MySQL 8.0 Reference Manual

Abstract

This is the MySQL™ Reference Manual. It documents MySQL 8.0 through 8.0.3.

MySQL 8.0 status. MySQL 8.0 is a development series not recommended for production use.

MySQL Cluster is currently not supported in MySQL 8.0. For information about MySQL Cluster, please see MySQL NDB Cluster 7.5 and NDB Cluster 7.6.

MySQL 8.0 features. This manual describes features that are not included in every edition of MySQL 8.0; such features may not be included in the edition of MySQL 8.0 licensed to you. If you have any questions about the features included in your edition of MySQL 8.0, refer to your MySQL 8.0 license agreement or contact your Oracle sales representative.

For notes detailing the changes in each release, see the MySQL 8.0 Release Notes.

For legal information, see the Legal Notices.

For help with using MySQL, please visit either the MySQL Forums or MySQL Mailing Lists, where you can discuss your issues with other MySQL users.

For additional documentation on MySQL products, including translations of the documentation into other languages, and downloadable versions in variety of formats, including HTML and PDF formats, see the MySQL Documentation Library.

Licensing information—MySQL 8.0. This product may include third-party software, used under license. If you are using a *Community* release of MySQL 8.0, see this document for licensing information, including licensing information relating to third-party software that may be included in this Community release.

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Preface and Legal Notices

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Chapter 1 General Information

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The MySQL™ software delivers a very fast, multi-threaded, multi-user, and robust SQL (Structured Query Language) database server. MySQL Server is intended for mission-critical, heavy-load production systems as well as for embedding into mass-deployed software. Oracle is a registered trademark of Oracle Corporation and/or its affiliates. MySQL is a trademark of Oracle Corporation and/or its affiliates, and shall not be used by Customer without Oracle's express written authorization. Other names may be trademarks of their respective owners.

The MySQL software is Dual Licensed. Users can choose to use the MySQL software as an Open Source product under the terms of the GNU General Public License (http://www.fsf.org/licenses/) or can purchase a standard commercial license from Oracle. See http://www.mysql.com/company/legal/licensing/ for more information on our licensing policies.

The following list describes some sections of particular interest in this manual:

- For a discussion of MySQL Database Server capabilities, see Section 1.3.2, “The Main Features of MySQL”.
- For an overview of new MySQL features, see Section 1.4, “What Is New in MySQL 8.0”. For information about the changes in each version, see the Release Notes.
- For installation instructions, see Chapter 2, *Installing and Upgrading MySQL*. For information about upgrading MySQL, see Section 2.10.1, “Upgrading MySQL”.
- For a tutorial introduction to the MySQL Database Server, see Chapter 3, *Tutorial*.
- For information about configuring and administering MySQL Server, see Chapter 5, *MySQL Server Administration*.

- For information about security in MySQL, see Chapter 6, *Security*.
- For information about setting up replication servers, see Chapter 18, *Replication*.
- For information about MySQL Enterprise, the commercial MySQL release with advanced features and management tools, see Chapter 29, *MySQL Enterprise Edition*.
- For answers to a number of questions that are often asked concerning the MySQL Database Server and its capabilities, see Appendix A, *MySQL 8.0 Frequently Asked Questions*.
- For a history of new features and bug fixes, see the Release Notes.

Important

To report problems or bugs, please use the instructions at Section 1.7, "How to Report Bugs or Problems". If you find a sensitive security bug in MySQL Server, please let us know immediately by sending an email message to `<secalert_us@oracle.com>`. Exception: Support customers should report all problems, including security bugs, to Oracle Support.

1.1 About This Manual

This is the Reference Manual for the MySQL Database System, version 8.0, through release 8.0.3. Differences between minor versions of MySQL 8.0 are noted in the present text with reference to release numbers (8.0.*x*). For license information, see the Legal Notices.

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Because this manual serves as a reference, it does not provide general instruction on SQL or relational database concepts. It also does not teach you how to use your operating system or command-line interpreter.

The MySQL Database Software is under constant development, and the Reference Manual is updated frequently as well. The most recent version of the manual is available online in searchable form at http://dev.mysql.com/doc/. Other formats also are available there, including HTML, PDF, and EPUB versions.

The Reference Manual source files are written in DocBook XML format. The HTML version and other formats are produced automatically, primarily using the DocBook XSL stylesheets. For information about DocBook, see http://docbook.org/

The source code for MySQL itself contains internal documentation written using Doxygen. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, "Generating MySQL Doxygen Documentation Content".

If you have questions about using MySQL, you can ask them using our mailing lists or forums. See Section 1.6.2, "MySQL Mailing Lists", and Section 1.6.3, "MySQL Community Support at the MySQL Forums". If you have suggestions concerning additions or corrections to the manual itself, please send them to the http://www.mysql.com/company/contact/.

This manual was originally written by David Axmark and Michael "Monty" Widenius. It is maintained by the MySQL Documentation Team, consisting of Chris Cole, Paul DuBois, Edward Gilmore, Stefan Hinz, David Moss, Philip Olson, Daniel Price, Daniel So, and Jon Stephens.

1.2 Typographical and Syntax Conventions

This manual uses certain typographical conventions:

- `Text in this style` is used for SQL statements; database, table, and column names; program listings and source code; and environment variables. Example: “To reload the grant tables, use the `FLUSH PRIVILEGES` statement.”
- **`Text in this style`** indicates input that you type in examples.
- `Text in this style` indicates the names of executable programs and scripts, examples being `mysql` (the MySQL command-line client program) and `mysqld` (the MySQL server executable).
- *`Text in this style`* is used for variable input for which you should substitute a value of your own choosing.
- *Text in this style* is used for emphasis.
- **Text in this style** is used in table headings and to convey especially strong emphasis.
- `Text in this style` is used to indicate a program option that affects how the program is executed, or that supplies information that is needed for the program to function in a certain way. *Example*: “The `--host` option (short form `-h`) tells the `mysql` client program the hostname or IP address of the MySQL server that it should connect to”.
- File names and directory names are written like this: “The global `my.cnf` file is located in the `/etc` directory.”
- Character sequences are written like this: “To specify a wildcard, use the ‘`%`’ character.”

When commands are shown that are meant to be executed from within a particular program, the prompt shown preceding the command indicates which command to use. For example, `shell>` indicates a command that you execute from your login shell, `root-shell>` is similar but should be executed as `root`, and `mysql>` indicates a statement that you execute from the `mysql` client program:

```
shell> type a shell command here
root-shell> type a shell command as root here
mysql> type a mysql statement here
```

In some areas different systems may be distinguished from each other to show that commands should be executed in two different environments. For example, while working with replication the commands might be prefixed with `master` and `slave`:

```
master> type a mysql command on the replication master here
slave> type a mysql command on the replication slave here
```

The “shell” is your command interpreter. On Unix, this is typically a program such as `sh`, `csh`, or `bash`. On Windows, the equivalent program is `command.com` or `cmd.exe`, typically run in a console window.

When you enter a command or statement shown in an example, do not type the prompt shown in the example.

Database, table, and column names must often be substituted into statements. To indicate that such substitution is necessary, this manual uses *`db_name`*, *`tbl_name`*, and *`col_name`*. For example, you might see a statement like this:

```
mysql> SELECT col_name FROM db_name.tbl_name;
```

This means that if you were to enter a similar statement, you would supply your own database, table, and column names, perhaps like this:

```
mysql> SELECT author_name FROM biblio_db.author_list;
```

SQL keywords are not case sensitive and may be written in any lettercase. This manual uses uppercase.

In syntax descriptions, square brackets (“[” and “]”) indicate optional words or clauses. For example, in the following statement, `IF EXISTS` is optional:

```
DROP TABLE [IF EXISTS] tbl_name
```

When a syntax element consists of a number of alternatives, the alternatives are separated by vertical bars (“|”). When one member from a set of choices *may* be chosen, the alternatives are listed within square brackets (“[” and “]”):

```
TRIM([[BOTH | LEADING | TRAILING] [remstr] FROM] str)
```

When one member from a set of choices *must* be chosen, the alternatives are listed within braces (“{” and “}”):

```
{DESCRIBE | DESC} tbl_name [col_name | wild]
```

An ellipsis (`...`) indicates the omission of a section of a statement, typically to provide a shorter version of more complex syntax. For example, `SELECT ... INTO OUTFILE` is shorthand for the form of `SELECT` statement that has an `INTO OUTFILE` clause following other parts of the statement.

An ellipsis can also indicate that the preceding syntax element of a statement may be repeated. In the following example, multiple *`reset_option`* values may be given, with each of those after the first preceded by commas:

```
RESET reset_option [,reset_option] ...
```

Commands for setting shell variables are shown using Bourne shell syntax. For example, the sequence to set the `CC` environment variable and run the `configure` command looks like this in Bourne shell syntax:

```
shell> CC=gcc ./configure
```

If you are using `csh` or `tcsh`, you must issue commands somewhat differently:

```
shell> setenv CC gcc
shell> ./configure
```

1.3 Overview of the MySQL Database Management System

1.3.1 What is MySQL?

MySQL, the most popular Open Source SQL database management system, is developed, distributed, and supported by Oracle Corporation.

The MySQL Web site (http://www.mysql.com/) provides the latest information about MySQL software.

- **MySQL is a database management system.**

  A database is a structured collection of data. It may be anything from a simple shopping list to a picture gallery or the vast amounts of information in a corporate network. To add, access, and process data stored in a computer database, you need a database management system such as MySQL Server. Since computers are very good at handling large amounts of data, database management systems play a central role in computing, as standalone utilities, or as parts of other applications.

- **MySQL databases are relational.**

  A relational database stores data in separate tables rather than putting all the data in one big storeroom. The database structures are organized into physical files optimized for speed. The

logical model, with objects such as databases, tables, views, rows, and columns, offers a flexible programming environment. You set up rules governing the relationships between different data fields, such as one-to-one, one-to-many, unique, required or optional, and “pointers” between different tables. The database enforces these rules, so that with a well-designed database, your application never sees inconsistent, duplicate, orphan, out-of-date, or missing data.

The SQL part of “MySQL” stands for “Structured Query Language”. SQL is the most common standardized language used to access databases. Depending on your programming environment, you might enter SQL directly (for example, to generate reports), embed SQL statements into code written in another language, or use a language-specific API that hides the SQL syntax.

SQL is defined by the ANSI/ISO SQL Standard. The SQL standard has been evolving since 1986 and several versions exist. In this manual, “SQL-92” refers to the standard released in 1992, “SQL:1999” refers to the standard released in 1999, and “SQL:2003” refers to the current version of the standard. We use the phrase “the SQL standard” to mean the current version of the SQL Standard at any time.

- **MySQL software is Open Source.**

  Open Source means that it is possible for anyone to use and modify the software. Anybody can download the MySQL software from the Internet and use it without paying anything. If you wish, you may study the source code and change it to suit your needs. The MySQL software uses the GPL (GNU General Public License), http://www.fsf.org/licenses/, to define what you may and may not do with the software in different situations. If you feel uncomfortable with the GPL or need to embed MySQL code into a commercial application, you can buy a commercially licensed version from us. See the MySQL Licensing Overview for more information (http://www.mysql.com/company/legal/licensing/).

- **The MySQL Database Server is very fast, reliable, scalable, and easy to use.**

  If that is what you are looking for, you should give it a try. MySQL Server can run comfortably on a desktop or laptop, alongside your other applications, web servers, and so on, requiring little or no attention. If you dedicate an entire machine to MySQL, you can adjust the settings to take advantage of all the memory, CPU power, and I/O capacity available. MySQL can also scale up to clusters of machines, networked together.

  MySQL Server was originally developed to handle large databases much faster than existing solutions and has been successfully used in highly demanding production environments for several years. Although under constant development, MySQL Server today offers a rich and useful set of functions. Its connectivity, speed, and security make MySQL Server highly suited for accessing databases on the Internet.

- **MySQL Server works in client/server or embedded systems.**

  The MySQL Database Software is a client/server system that consists of a multi-threaded SQL server that supports different back ends, several different client programs and libraries, administrative tools, and a wide range of application programming interfaces (APIs).

  We also provide MySQL Server as an embedded multi-threaded library that you can link into your application to get a smaller, faster, easier-to-manage standalone product.

- **A large amount of contributed MySQL software is available.**

  MySQL Server has a practical set of features developed in close cooperation with our users. It is very likely that your favorite application or language supports the MySQL Database Server.

The official way to pronounce “MySQL” is “My Ess Que Ell” (not “my sequel”), but we do not mind if you pronounce it as “my sequel” or in some other localized way.

1.3.2 The Main Features of MySQL

This section describes some of the important characteristics of the MySQL Database Software. In most respects, the roadmap applies to all versions of MySQL. For information about features as they are introduced into MySQL on a series-specific basis, see the “In a Nutshell” section of the appropriate Manual:

- MySQL 8.0: Section 1.4, “What Is New in MySQL 8.0”
- MySQL 5.7: What Is New in MySQL 5.7
- MySQL 5.6: What Is New in MySQL 5.6
- MySQL 5.5: What Is New in MySQL 5.5

Internals and Portability

- Written in C and C++.
- Tested with a broad range of different compilers.
- Works on many different platforms. See http://www.mysql.com/support/supportedplatforms/database.html.
- For portability, uses `CMake` in MySQL 5.5 and up. Previous series use GNU Automake, Autoconf, and Libtool.
- Tested with Purify (a commercial memory leakage detector) as well as with Valgrind, a GPL tool (http://developer.kde.org/~sewardj/).
- Uses multi-layered server design with independent modules.
- Designed to be fully multi-threaded using kernel threads, to easily use multiple CPUs if they are available.
- Provides transactional and nontransactional storage engines.
- Uses very fast B-tree disk tables (`MyISAM`) with index compression.
- Designed to make it relatively easy to add other storage engines. This is useful if you want to provide an SQL interface for an in-house database.
- Uses a very fast thread-based memory allocation system.
- Executes very fast joins using an optimized nested-loop join.
- Implements in-memory hash tables, which are used as temporary tables.
- Implements SQL functions using a highly optimized class library that should be as fast as possible. Usually there is no memory allocation at all after query initialization.
- Provides the server as a separate program for use in a client/server networked environment, and as a library that can be embedded (linked) into standalone applications. Such applications can be used in isolation or in environments where no network is available.

Data Types

- Many data types: signed/unsigned integers 1, 2, 3, 4, and 8 bytes long, `FLOAT`, `DOUBLE`, `CHAR`, `VARCHAR`, `BINARY`, `VARBINARY`, `TEXT`, `BLOB`, `DATE`, `TIME`, `DATETIME`, `TIMESTAMP`, `YEAR`, `SET`, `ENUM`, and OpenGIS spatial types. See Chapter 11, *Data Types*.
- Fixed-length and variable-length string types.

Statements and Functions

- Full operator and function support in the `SELECT` list and `WHERE` clause of queries. For example:

```
mysql> SELECT CONCAT(first_name, ' ', last_name)
    -> FROM citizen
    -> WHERE income/dependents > 10000 AND age > 30;
```

- Full support for SQL `GROUP BY` and `ORDER BY` clauses. Support for group functions (`COUNT()`, `AVG()`, `STD()`, `SUM()`, `MAX()`, `MIN()`, and `GROUP_CONCAT()`).
- Support for `LEFT OUTER JOIN` and `RIGHT OUTER JOIN` with both standard SQL and ODBC syntax.
- Support for aliases on tables and columns as required by standard SQL.
- Support for `DELETE`, `INSERT`, `REPLACE`, and `UPDATE` to return the number of rows that were changed (affected), or to return the number of rows matched instead by setting a flag when connecting to the server.
- Support for MySQL-specific `SHOW` statements that retrieve information about databases, storage engines, tables, and indexes. Support for the `INFORMATION_SCHEMA` database, implemented according to standard SQL.
- An `EXPLAIN` statement to show how the optimizer resolves a query.
- Independence of function names from table or column names. For example, `ABS` is a valid column name. The only restriction is that for a function call, no spaces are permitted between the function name and the "`(`" that follows it. See Section 9.3, "Keywords and Reserved Words".
- You can refer to tables from different databases in the same statement.

Security

- A privilege and password system that is very flexible and secure, and that enables host-based verification.
- Password security by encryption of all password traffic when you connect to a server.

Scalability and Limits

- Support for large databases. We use MySQL Server with databases that contain 50 million records. We also know of users who use MySQL Server with 200,000 tables and about 5,000,000,000 rows.
- Support for up to 64 indexes per table. Each index may consist of 1 to 16 columns or parts of columns. The maximum index width for `InnoDB` tables is either 767 bytes or 3072 bytes. See Section 15.8.1.7, "Limits on InnoDB Tables". The maximum index width for `MyISAM` tables is 1000 bytes. See Section 16.2, "The MyISAM Storage Engine". An index may use a prefix of a column for `CHAR`, `VARCHAR`, `BLOB`, or `TEXT` column types.

Connectivity

- Clients can connect to MySQL Server using several protocols:
  - Clients can connect using TCP/IP sockets on any platform.
  - On Windows systems, clients can connect using named pipes if the server is started with the `--enable-named-pipe` option. Windows servers also support shared-memory connections if started with the `--shared-memory` option. Clients can connect through shared memory by using the `--protocol=memory` option.
  - On Unix systems, clients can connect using Unix domain socket files.
- MySQL client programs can be written in many languages. A client library written in C is available for clients written in C or C++, or for any language that provides C bindings.
- APIs for C, C++, Eiffel, Java, Perl, PHP, Python, Ruby, and Tcl are available, enabling MySQL clients to be written in many languages. See Chapter 27, *Connectors and APIs*.

- The Connector/ODBC (MyODBC) interface provides MySQL support for client programs that use ODBC (Open Database Connectivity) connections. For example, you can use MS Access to connect to your MySQL server. Clients can be run on Windows or Unix. Connector/ODBC source is available. All ODBC 2.5 functions are supported, as are many others. See MySQL Connector/ODBC Developer Guide.
- The Connector/J interface provides MySQL support for Java client programs that use JDBC connections. Clients can be run on Windows or Unix. Connector/J source is available. See MySQL Connector/J 5.1 Developer Guide.
- MySQL Connector/Net enables developers to easily create .NET applications that require secure, high-performance data connectivity with MySQL. It implements the required ADO.NET interfaces and integrates into ADO.NET aware tools. Developers can build applications using their choice of .NET languages. MySQL Connector/Net is a fully managed ADO.NET driver written in 100% pure C#. See MySQL Connector/Net Developer Guide.

Localization

- The server can provide error messages to clients in many languages. See Section 10.2, “Setting the Error Message Language”.
- Full support for several different character sets, including `latin1` (cp1252), `german`, `big5`, `ujis`, several Unicode character sets, and more. For example, the Scandinavian characters “å”, “ä” and “ö” are permitted in table and column names.
- All data is saved in the chosen character set.
- Sorting and comparisons are done according to the default character set and collation. is possible to change this when the MySQL server is started (see Section 10.1.3.2, “Server Character Set and Collation”). To see an example of very advanced sorting, look at the Czech sorting code. MySQL Server supports many different character sets that can be specified at compile time and runtime.
- The server time zone can be changed dynamically, and individual clients can specify their own time zone. See Section 10.6, “MySQL Server Time Zone Support”.

Clients and Tools

- MySQL includes several client and utility programs. These include both command-line programs such as `mysqldump` and `mysqladmin`, and graphical programs such as MySQL Workbench.
- MySQL Server has built-in support for SQL statements to check, optimize, and repair tables. These statements are available from the command line through the `mysqlcheck` client. MySQL also includes `myisamchk`, a very fast command-line utility for performing these operations on `MyISAM` tables. See Chapter 4, *MySQL Programs*.
- MySQL programs can be invoked with the `--help` or `-?` option to obtain online assistance.

1.3.3 History of MySQL

We started out with the intention of using the `mSQL` database system to connect to our tables using our own fast low-level (ISAM) routines. However, after some testing, we came to the conclusion that `mSQL` was not fast enough or flexible enough for our needs. This resulted in a new SQL interface to our database but with almost the same API interface as `mSQL`. This API was designed to enable third-party code that was written for use with `mSQL` to be ported easily for use with MySQL.

MySQL is named after co-founder Monty Widenius's daughter, My.

The name of the MySQL Dolphin (our logo) is “Sakila,” which was chosen from a huge list of names suggested by users in our “Name the Dolphin” contest. The winning name was submitted by Ambrose Twebaze, an Open Source software developer from Swaziland, Africa. According to Ambrose, the feminine name Sakila has its roots in SiSwati, the local language of Swaziland. Sakila is also the name of a town in Arusha, Tanzania, near Ambrose's country of origin, Uganda.

1.4 What Is New in MySQL 8.0

This section summarizes what has been added to, deprecated in, and removed from MySQL 8.0. A companion section lists MySQL server options and variables that have been added, deprecated, or removed in MySQL 8.0. See Section 1.5, “Server and Status Variables and Options Added, Deprecated, or Removed in MySQL 8.0”.

- Features Added in MySQL 8.0
- Features Deprecated in MySQL 8.0
- Features Removed in MySQL 8.0

Features Added in MySQL 8.0

The following features have been added to MySQL 8.0:

- **Data dictionary.** These enhancements were added:
  - MySQL Server now incorporates a transactional data dictionary that stores information about database objects. In previous MySQL releases, dictionary data was stored in metadata files and nontransactional tables.

    `InnoDB` continues to use its own data dictionary in the MySQL 8.0.0 release.

    For more information about this feature, see Chapter 14, *MySQL Data Dictionary*.
- **Account management.** MySQL now supports roles, which are named collections of privileges. Roles can be created and dropped. Roles can have privileges granted to and revoked from them. Roles can be granted to and revoked from user accounts. The active applicable roles for an account can be selected from among those granted to the account, and can be changed during sessions for that account. For more information, see Section 6.3.4, “Using Roles”.
- **InnoDB enhancements.** These `InnoDB` enhancements were added:
  - The current maximum auto-increment counter value is written to the redo log each time the value changes, and it is saved to an engine-private system table on each checkpoint. These changes make the current maximum auto-increment counter value persistent across server restarts. Additionally:
    - A server restart no longer cancels the effect of the `AUTO_INCREMENT = N` table option. If you initialize the auto-increment counter to a specific value, or if you alter the auto-increment counter value to a larger value, the new value is persisted across server restarts.
    - A server restart immediately following a `ROLLBACK` operation no longer results in the reuse of auto-increment values that were allocated to the rolled-back transaction.
    - If you modify an `AUTO_INCREMENT` column value to a value larger than the current maximum auto-increment value (in an `UPDATE` operation, for example), the new value is persisted, and subsequent `INSERT` operations allocate auto-increment values starting from the new, larger value.

    For more information, see Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”, and InnoDB AUTO_INCREMENT Counter Initialization.
  - When encountering index tree corruption, `InnoDB` writes a corruption flag to the redo log, which makes the corruption flag crash-safe. `InnoDB` also writes in-memory corruption flag data to an engine-private system table on each checkpoint. During recovery, `InnoDB` reads corruption flags from both locations and merges results before marking in-memory table and index objects as corrupt.

- The InnoDB memcached plugin supports multiple get operations (fetching multiple key/value pairs in a single memcached query) and range queries. See Section 15.19.4, “InnoDB memcached Multiple get and Range Query Support”.

- A new dynamic configuration option, innodb_deadlock_detect, is used to disable deadlock detection. On high concurrency systems, deadlock detection can cause a slowdown when numerous threads wait for the same lock. At times, it may be more efficient to disable deadlock detection and rely on the innodb_lock_wait_timeout setting for transaction rollback when a deadlock occurs.

- A new INFORMATION_SCHEMA table, INNODB_CACHED_INDEXES, reports the number of index pages cached in the InnoDB buffer pool for each index.

- All InnoDB temporary tables are created in the shared temporary tablespace, ibtmp1.

- The InnoDB tablespace encryption feature supports encryption of redo log and undo log data. See Redo Log Data Encryption, and Undo Log Data Encryption.

- InnoDB supports NOWAIT and SKIP LOCKED options with SELECT ... FOR SHARE and SELECT ... FOR UPDATE locking read statements. NOWAIT causes the statement to return immediately if a requested row is locked by another transaction. SKIP LOCKED removes locked rows from the result set. See Locking Read Concurrency with NOWAIT and SKIP LOCKED.

  SELECT ... FOR SHARE replaces SELECT ... LOCK IN SHARE MODE, but LOCK IN SHARE MODE remains available for backward compatibility. The statements are equivalent. However, FOR UPDATE and FOR SHARE support NOWAIT, SKIP LOCKED, and OF *tbl_name* options. See Section 13.2.10, “SELECT Syntax”.

  OF *tbl_name* applies locking queries to named tables.

- As of 8.0.2, ADD PARTITION, DROP PARTITION, COALESCE PARTITION, REORGANIZE PARTITION, and REBUILD PARTITION ALTER TABLE options are supported by native partitioning in-place APIs and may be used with ALGORITHM={COPY|INPLACE} and LOCK clauses.

  DROP PARTITION with ALGORITHM=INPLACE deletes data stored in the partition and drops the partition. However, DROP PARTITION with ALGORITHM=COPY or old_alter_table=ON rebuilds the partitioned table and attempts to move data from the dropped partition to another partition with a compatible PARTITION ... VALUES definition. Data that cannot be moved to another partition is deleted.

- As of MySQL 8.0.2, the InnoDB storage engine uses the global MySQL data dictionary rather than its own storage engine-specific data dictionary. For information about the data dictionary, see Chapter 14, *MySQL Data Dictionary*.

- As of MySQL 8.0.2, mysql system tables and data dictionary tables are created in a single InnoDB tablespace file named mysql.ibd in the MySQL data directory. Previously, these tables were created in individual InnoDB tablespace files in the mysql database directory.

- **JSON enhancements.** The following enhancements or additions were made to MySQL's JSON functionality:

  - Added the ->> (inline path) operator, which is equivalent to calling JSON_UNQUOTE() on the result of JSON_EXTRACT().

    This is a refinement of the column path operator -> introduced in MySQL 5.7; col->>"$.path" is equivalent to JSON_UNQUOTE(col->"$.path"). The inline path operator can be used wherever you can use JSON_UNQUOTE(JSON_EXTRACT()), such SELECT column lists, WHERE

and `HAVING` clauses, and `ORDER BY` and `GROUP BY` clauses. For more information, see the description of the operator, as well as Section 12.16.7, “JSON Path Syntax”.

- Added two JSON aggregation functions `JSON_ARRAYAGG()` and `JSON_OBJECTAGG()`. `JSON_ARRAYAGG()` takes a column or expression as its argument, and aggregates the result as a single `JSON` array. The expression can evaluate to any MySQL data type; this does not have to be a `JSON` value. `JSON_OBJECTAGG()` takes two columns or expressions which it interprets as a key and a value; it returns the result as a single `JSON` object. For more information and examples, see Section 12.18, “Aggregate (GROUP BY) Functions”.

- Added the JSON utility function `JSON_PRETTY()`, which outputs an existing `JSON` value in an easy-to-read format; each JSON object member or array value is printed on a separate line, and a child object or array is intended 2 spaces with respect to its parent.

  This function also works with a string that can be parsed as a JSON value.

  For more detailed information and examples, see Section 12.16.6, “JSON Utility Functions”.

- When sorting `JSON` values in a query using `ORDER BY`, each value is now represented by a variable-length part of the sort key, rather than a part of a fixed 1K in size. In many cases this can reduce excessive usage; for example, a scalar `INT` or even `BIGINT` value actually requires very few bytes, so that the remainder of this space (up to 90% or more) was taken up by padding. This change has the following benefits for performance:

  - Sort buffer space is now used more effectively, so that filesorts need not flush to disk as early or often as with fixed-length sort keys. This means that more data can be sorted in memory, avoiding unnecessary disk access.

  - Shorter keys can be compared more quickly than longer ones, providing a noticeable improvement in performance. This is true for sorts performed entirely in memory as well as for sorts that require writing to and reading from disk.

- Added support in MySQL 8.0.2 for partial, in-place updates of `JSON` column values, which is more efficient than completely removing an existing JSON value and writing a new one in its place, as was done previously when updating any `JSON` column. For this optimization to be applied, the update must be applied using `JSON_SET()`, `JSON_REPLACE()`, or `JSON_REMOVE()`. New elements cannot be added to the JSON document being updated; values within the document cannot take more space than previous to the update. See Section 11.6, “The JSON Data Type”, for a detailed discussion of the requirements.

- Added the JSON utility functions `JSON_STORAGE_SIZE()` and `JSON_STORAGE_FREE()`. `JSON_STORAGE_SIZE()` returns the storage space in bytes used for the binary representation of a JSON document prior to any partial update (see previous item). `JSON_STORAGE_FREE()` shows the amount of space remaining in a table column of type `JSON` after it has been partially updated using `JSON_SET()` or `JSON_REPLACE()`; this is greater than zero if the binary representation of the new value is less than that of the previous value.

  Each of these functions also accepts a valid string representation of a JSON document. For such a value, `JSON_STORAGE_SIZE()` returns the space used by its binary representation following its conversion to a JSON document. For a variable containing the string representation of a JSON document, `JSON_STORAGE_FREE()` returns zero. Either function produces an error if its (non-null) argument cannot be parsed as a valid JSON document, and `NULL` if the argument is `NULL`.

  For more information and examples, see Section 12.16.6, “JSON Utility Functions”.

  `JSON_STORAGE_SIZE()` and `JSON_STORAGE_FREE()` were implemented in MySQL 8.0.2.

- Added support in MySQL 8.0.2 for ranges such as `$[1 to 5]` in XPath expressions. Also added support in this version for the `last` keyword and relative addressing, such that `$[last]` always selects the last (highest-numbered) element in the array and `$[last-1]` the next to last element.

`last` and expressions using it can also be included in range definitions; for example, `$[last-2 to last-1]` returns the last two elements but one from an array. See Searching and Modifying JSON Values, for additional information and examples.

- **Optimizer.** These optimizer enhancements were added:
  - MySQL now supports invisible indexes. An invisible index is not used by the optimizer at all, but is otherwise maintained normally. Indexes are visible by default. Invisible indexes make it possible to test the effect of removing an index on query performance, without making a destructive change that must be undone should the index turn out to be required. See Section 8.3.11, “Invisible Indexes”.
  - MySQL now supports descending indexes: `DESC` in an index definition is no longer ignored but causes storage of key values in descending order. Previously, indexes could be scanned in reverse order but at a performance penalty. A descending index can be scanned in forward order, which is more efficient. Descending indexes also make it possible for the optimizer to use multiple-column indexes when the most efficient scan order mixes ascending order for some columns and descending order for others. See Section 8.3.12, “Descending Indexes”.
- **Common table expressions.** MySQL now supports common table expressions, both nonrecursive and recursive. Common table expressions enable use of named temporary result sets, implemented by permitting a `WITH` clause preceding `SELECT` statements and certain other statements. For more information, see Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”.
- **Internal Temporary Tables.** As of MySQL 8.0.2, the `TempTable` storage engine replaces the `MEMORY` storage engine as the default engine for in-memory internal temporary tables. The `TempTable` storage engine provides efficient storage for `VARCHAR` and `VARBINARY` columns. The `internal_tmp_mem_storage_engine` session variable defines the storage engine for in-memory internal temporary tables. Permitted values are `TempTable` (the default) and `MEMORY`. The `temptable_max_ram` configuration option defines the maximum amount of memory that the `TempTable` storage engine can use before data is stored to disk.

Features Deprecated in MySQL 8.0

The following features are deprecated in MySQL 8.0 and may be or will be removed in a future series. Where alternatives are shown, applications should be updated to use them.

Features Removed in MySQL 8.0

The following items are obsolete and have been removed in MySQL 8.0. Where alternatives are shown, applications should be updated to use them.

- The data dictionary provides information about database objects, so the server no longer checks directory names in the data directory to find databases. Consequently, the `--ignore-db-dir` option and `ignore_db_dirs` system variables are extraneous and have been removed.
- The `sync_frm` system variable has been removed because `.frm` files have become obsolete.
- In MySQL 5.7, several spatial functions available under multiple names were deprecated to move in the direction of making the spatial function namespace more consistent, the goal being that each spatial function name begin with `ST_` if it performs an exact operation, or with `MBR` if it performs an operation based on minimum bounding rectangles. In MySQL 8.0, the deprecated functions are removed to leave only the corresponding `ST_` and `MBR` functions:
  - These functions are removed in favor of the `MBR` names: `Contains()`, `Disjoint()`, `Equals()`, `Intersects()`, `Overlaps()`, `Within()`.
  - These functions are removed in favor of the `ST_` names: `Area()`, `AsBinary()`, `AsText()`, `AsWKB()`, `AsWKT()`, `Buffer()`, `Centroid()`, `ConvexHull()`, `Crosses()`,

`Dimension()`, `Distance()`, `EndPoint()`, `Envelope()`, `ExteriorRing()`, `GeomCollFromText()`, `GeomCollFromWKB()`, `GeomFromText()`, `GeomFromWKB()`, `GeometryCollectionFromText()`, `GeometryCollectionFromWKB()`, `GeometryFromText()`, `GeometryFromWKB()`, `GeometryN()`, `GeometryType()`, `InteriorRingN()`, `IsClosed()`, `IsEmpty()`, `IsSimple()`, `LineFromText()`, `LineFromWKB()`, `LineStringFromText()`, `LineStringFromWKB()`, `MLineFromText()`, `MLineFromWKB()`, `MPointFromText()`, `MPointFromWKB()`, `MPolyFromText()`, `MPolyFromWKB()`, `MultiLineStringFromText()`, `MultiLineStringFromWKB()`, `MultiPointFromText()`, `MultiPointFromWKB()`, `MultiPolygonFromText()`, `MultiPolygonFromWKB()`, `NumGeometries()`, `NumInteriorRings()`, `NumPoints()`, `PointFromText()`, `PointFromWKB()`, `PointN()`, `PolyFromText()`, `PolyFromWKB()`, `PolygonFromText()`, `PolygonFromWKB()`, `SRID()`, `StartPoint()`, `Touches()`, `X()`, `Y()`.

  - `GLength()` is removed in favor of `ST_Length()`.

- The functions described in Section 12.15.4, “Functions That Create Geometry Values from WKB Values” previously accepted either WKB strings or geometry arguments. Geometry arguments are no longer permitted and produce an error. See that section for guidelines for migrating queries away from using geometry arguments.

- The parser no longer treats `\N` as a synonym for `NULL` in SQL statements. Use `NULL` instead.

  This change does not affect text file import or export operations performed with `LOAD DATA INFILE` or `SELECT ... INTO OUTFILE`, for which `NULL` continues to be represented by `\N`. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

- `PROCEDURE ANALYSE()` syntax is removed.

- The client-side `--ssl` and `--ssl-verify-server-cert` options have been removed. Use `--ssl-mode=REQUIRED` instead of `--ssl=1` or `--enable-ssl`. Use `--ssl-mode=DISABLED` instead of `--ssl=0`, `--skip-ssl`, or `--disable-ssl`. Use `--ssl-mode=VERIFY_IDENTITY` instead of `--ssl-verify-server-cert` options. (The server-side `--ssl` option remains unchanged.)

  For the C API, `MYSQL_OPT_SSL_ENFORCE` and `MYSQL_OPT_SSL_VERIFY_SERVER_CERT` options for `mysql_options()` correspond to the client-side `--ssl` and `--ssl-verify-server-cert` options and have been removed. Use `MYSQL_OPT_SSL_MODE` with an option value of `SSL_MODE_REQUIRED` or `SSL_MODE_VERIFY_IDENTITY` instead.

- The `--temp-pool` server option has been removed.

- The server no longer performs conversion of pre-MySQL 5.1 database names containing special characters to 5.1 format with the addition of a `#mysql50#` prefix. Because these conversions are no longer performed, the `--fix-db-names` and `--fix-table-names` options for `mysqlcheck`, the `UPGRADE DATA DIRECTORY NAME` clause for the `ALTER DATABASE` statement, and the `Com_alter_db_upgrade` status variable have been removed.

  Upgrades are supported only from one major version to another (for example, 5.0 to 5.1, or 5.1 to 5.5), so there should be little remaining need for conversion of older 5.0 database names to current versions of MySQL. As a workaround, upgrade a MySQL 5.0 installation to MySQL 5.1 before upgrading to a more recent release.

- The `mysql_install_db` program has been removed from MySQL distributions. Data directory initialization should be performed by invoking `mysqld` with the `--initialize` or `--initialize-insecure` option instead. In addition, the `--bootstrap` option for `mysqld` that was used by `mysql_install_db` has been removed, and the `INSTALL_SCRIPTDIR` `CMake` option that controlled the installation location for `mysql_install_db` has been removed.

- The generic partitioning handler has been removed from the MySQL server. In order to support partitioning of a given table, the storage engine used for the table must now provide its own (“native”) partitioning handler. The `--partition` and `--skip-partition` options have been removed

from the MySQL Server, and partitioning-related entries are no longer shown in the output of `SHOW PLUGINS` or in the `INFORMATION_SCHEMA.PLUGINS` table.

Two MySQL storage engines currently provide native partitioning support—`InnoDB` and `NDB`; of these, only `InnoDB` is supported in MySQL 8.0. Any attempt to create partitioned tables in MySQL 8.0 using any other storage engine fails.

Ramifications for upgrades. The direct upgrade of a partitioned table using a storage engine other than `InnoDB` (such as `MyISAM`) from MySQL 5.7 (or earlier) to MySQL 8.0 is not supported. There are two options for handling such a table:

- Remove the table's partitioning, using `ALTER TABLE ... REMOVE PARTITIONING`.
- Change the storage engine used for the table to `InnoDB`, with `ALTER TABLE ... ENGINE=INNODB`.

At least one of the two operations just listed must be performed for each partitioned non-`InnoDB` table prior to upgrading the server to MySQL 8.0. Otherwise, such a table cannot be used following the upgrade.

Due to the fact that table creation statements that would result in a partitioned table using a storage engine without partitioning support now fail with an error (`ER_CHECK_NOT_IMPLEMENTED`), you must make sure that any statements in a dump file (such as that written by `mysqldump`) from an older version of MySQL that you wish to import into a MySQL 8.0 server that create partitioned tables do not also specify a storage engine such as `MyISAM` that has no native partitioning handler. You can do this by performing either of the following:

- Remove any references to partitioning from `CREATE TABLE` statements that use a value for the `STORAGE ENGINE` option other than `InnoDB`.
- Specifying the storage engine as `InnoDB`, or allow `InnoDB` to be used as the table's storage engine by default.

For more information, see Section 22.6.2, “Partitioning Limitations Relating to Storage Engines”.

- System and status variable information is no longer maintained in the `INFORMATION_SCHEMA`. These tables have been removed: `GLOBAL_VARIABLES`, `SESSION_VARIABLES`, `GLOBAL_STATUS`, `SESSION_STATUS`. Use the corresponding Performance Schema tables instead. See Section 25.11.13, “Performance Schema System Variable Tables”, and Section 25.11.14, “Performance Schema Status Variable Tables”. In addition, the `show_compatibility_56` system variable has been removed. It was used in the transition period during which system and status variable information in `INFORMATION_SCHEMA` tables was moved to Performance Schema tables, and is no longer needed. These status variables have been removed: `Slave_heartbeat_period`, `Slave_last_heartbeat`, `Slave_received_heartbeats`, `Slave_retried_transactions`, `Slave_running`. The information they provided is available in Performance Schema tables; see Migrating to Performance Schema System and Status Variable Tables.

- The `libmysqld` embedded server library has been removed, along with:
  - The `mysql_config --libmysqld-libs`, `--embedded-libs`, and `--embedded` options
  - The `CMake WITH_EMBEDDED_SERVER`, `WITH_EMBEDDED_SHARED_LIBRARY`, and `INSTALL_SECURE_FILE_PRIV_EMBEDDEDDIR` options
  - The (undocumented) `mysql --server-arg` option
  - The `mysqltest --embedded-server`, `--server-arg`, and `--server-file` options
  - The `mysqltest_embedded` and `mysql_client_test_embedded` test programs

- The `mysql_plugin` utility has been removed. Alternatives include loading plugins at server startup using the `--plugin-load` or `--plugin-load-add` option, or at runtime using the `INSTALL PLUGIN` statement.

- The following server error codes are not used and have been removed. Applications that test specifically for any of these errors should be updated.

```
ER_BINLOG_READ_EVENT_CHECKSUM_FAILURE
ER_BINLOG_ROW_RBR_TO_SBR
ER_BINLOG_ROW_WRONG_TABLE_DEF
ER_CANT_ACTIVATE_LOG
ER_CANT_CHANGE_GTID_NEXT_IN_TRANSACTION
ER_CANT_CREATE_FEDERATED_TABLE
ER_CANT_CREATE_SROUTINE
ER_CANT_DELETE_FILE
ER_CANT_GET_WD
ER_CANT_SET_GTID_PURGED_WHEN_GTID_MODE_IS_OFF
ER_CANT_SET_WD
ER_CANT_WRITE_LOCK_LOG_TABLE
ER_CREATE_DB_WITH_READ_LOCK
ER_CYCLIC_REFERENCE
ER_DB_DROP_DELETE
ER_DELAYED_NOT_SUPPORTED
ER_DIFF_GROUPS_PROC
ER_DISK_FULL
ER_DROP_DB_WITH_READ_LOCK
ER_DROP_USER
ER_DUMP_NOT_IMPLEMENTED
ER_ERROR_DURING_CHECKPOINT
ER_ERROR_ON_CLOSE
ER_EVENTS_DB_ERROR
ER_EVENT_CANNOT_DELETE
ER_EVENT_CANT_ALTER
ER_EVENT_COMPILE_ERROR
ER_EVENT_DATA_TOO_LONG
ER_EVENT_DROP_FAILED
ER_EVENT_MODIFY_QUEUE_ERROR
ER_EVENT_NEITHER_M_EXPR_NOR_M_AT
ER_EVENT_OPEN_TABLE_FAILED
ER_EVENT_STORE_FAILED
ER_EXEC_STMT_WITH_OPEN_CURSOR
ER_FAILED_ROUTINE_BREAK_BINLOG
ER_FLUSH_MASTER_BINLOG_CLOSED
ER_FORM_NOT_FOUND
ER_FOUND_GTID_EVENT_WHEN_GTID_MODE_IS_OFF__UNUSED
ER_FRM_UNKNOWN_TYPE
ER_GOT_SIGNAL
ER_GRANT_PLUGIN_USER_EXISTS
ER_GTID_MODE_REQUIRES_BINLOG
ER_GTID_NEXT_IS_NOT_IN_GTID_NEXT_LIST
ER_HASHCHK
ER_INDEX_REBUILD
ER_INNODB_NO_FT_USES_PARSER
ER_LIST_OF_FIELDS_ONLY_IN_HASH_ERROR
ER_LOAD_DATA_INVALID_COLUMN_UNUSED
ER_LOGGING_PROHIBIT_CHANGING_OF
ER_MALFORMED_DEFINER
ER_MASTER_KEY_ROTATION_ERROR_BY_SE
ER_NDB_CANT_SWITCH_BINLOG_FORMAT
ER_NEVER_USED
ER_NISAMCHK
ER_NO_CONST_EXPR_IN_RANGE_OR_LIST_ERROR
ER_NO_FILE_MAPPING
ER_NO_GROUP_FOR_PROC
ER_NO_RAID_COMPILED
ER_NO_SUCH_KEY_VALUE
ER_NO_SUCH_PARTITION__UNUSED
ER_OBSOLETE_CANNOT_LOAD_FROM_TABLE
ER_OBSOLETE_COL_COUNT_DOESNT_MATCH_CORRUPTED
ER_ORDER_WITH_PROC
```

```
ER_PARTITION_SUBPARTITION_ERROR
ER_PARTITION_SUBPART_MIX_ERROR
ER_PART_STATE_ERROR
ER_PASSWD_LENGTH
ER_QUERY_ON_MASTER
ER_RBR_NOT_AVAILABLE
ER_SKIPPING_LOGGED_TRANSACTION
ER_SLAVE_CHANNEL_DELETE
ER_SLAVE_MULTIPLE_CHANNELS_HOST_PORT
ER_SLAVE_MUST_STOP
ER_SLAVE_WAS_NOT_RUNNING
ER_SLAVE_WAS_RUNNING
ER_SP_GOTO_IN_HNDLR
ER_SP_PROC_TABLE_CORRUPT
ER_SQL_MODE_NO_EFFECT
ER_SR_INVALID_CREATION_CTX
ER_TABLE_NEEDS_UPG_PART
ER_TOO_MUCH_AUTO_TIMESTAMP_COLS
ER_UNEXPECTED_EOF
ER_UNION_TABLES_IN_DIFFERENT_DIR
ER_UNSUPPORTED_BY_REPLICATION_THREAD
ER_UNUSED1
ER_UNUSED2
ER_UNUSED3
ER_UNUSED4
ER_UNUSED5
ER_UNUSED6
ER_VIEW_SELECT_DERIVED_UNUSED
ER_WRONG_MAGIC
ER_WSAS_FAILED
```

- The deprecated `INFORMATION_SCHEMA` `INNODB_LOCKS` and `INNODB_LOCK_WAITS` tables have been removed. Use the Performance Schema `data_locks` and `data_lock_waits` tables instead.

  Note

  In MySQL 5.7, the `LOCK_TABLE` column in the `INNODB_LOCKS` table and the `locked_table` column in the `sys` schema `innodb_lock_waits` and `x$innodb_lock_waits` views contain combined schema/table name values. In MySQL 8.0, the `data_locks` table and the `sys` schema views contain separate schema name and table name columns. See Section 26.4.3.9, “The innodb_lock_waits and x$innodb_lock_waits Views”.

- `InnoDB` no longer supports compressed temporary tables. When `innodb_strict_mode` is enabled (the default), `CREATE TEMPORARY TABLE` returns an error if `ROW_FORMAT=COMPRESSED` or `KEY_BLOCK_SIZE` is specified. If `innodb_strict_mode` is disabled, warnings are issued and the temporary table is created using a non-compressed row format.

- `InnoDB` no longer creates `.isl` files (`InnoDB` Symbolic Link files) when creating tablespace data files outside of the MySQL data directory. Redo log records are now used to locate remote tablespaces.

  With this change, moving a remote tablespace while the server is offline by manually modifying an `.isl` file is no longer supported.

- The following `InnoDB` file format configuration parameters were removed:

  - `innodb_file_format`
  - `innodb_file_format_check`
  - `innodb_file_format_max`
  - `innodb_large_prefix`

File format configuration parameters were necessary for creating tables compatible with earlier versions of `InnoDB` in MySQL 5.1. Now that MySQL 5.1 has reached the end of its product lifecycle, the parameters are no longer required.

The `FILE_FORMAT` column was removed from the `INNODB_SYS_TABLES` and `INNODB_SYS_TABLESPACES` Information Schema tables.

- The `innodb_support_xa` system variable, which enables support for two-phase commit in XA transactions, was removed. `InnoDB` support for two-phase commit in XA transactions is always enabled.
- Support for DTrace has been removed.

1.5 Server and Status Variables and Options Added, Deprecated, or Removed in MySQL 8.0

This section lists server variables, status variables, and options that were added for the first time, have been deprecated, or have been removed in MySQL 8.0.

- Options and Variables Introduced in MySQL 8.0
- Options and Variables Deprecated in MySQL 8.0
- Options and Variables Removed in MySQL 8.0

Where applicable, separate lists have been provided—for variables and options which have been added, removed, or deprecated —within each section.

Options and Variables Introduced in MySQL 8.0

The following system variables, status variables, and options are new in MySQL 8.0, and have not been included in any previous release series.

- `group_replication_member_weight`: Chance of this member being elected as primary. Added in MySQL 8.0.2.
- `group_replication_flow_control_hold_percent`: Defines what percentage of the group quota remains unused. Added in MySQL 8.0.2.
- `group_replication_flow_control_max_commit_quota`: Defines the maximum flow control quota of the group. Added in MySQL 8.0.2.
- `group_replication_flow_control_member_quota_percent`: Defines the percentage of the quota that a member should assume is available for itself when calculating the quotas. Added in MySQL 8.0.2.
- `group_replication_flow_control_min_quota`: Controls the lowest flow control quota that can be assigned to a member. Added in MySQL 8.0.2.
- `group_replication_flow_control_min_recovery_quota`: Controls the lowest quota that can be assigned to a member because of another recovering member in the group. Added in MySQL 8.0.2.
- `group_replication_flow_control_period`: Defines how many seconds to wait between flow control iterations. Added in MySQL 8.0.2.
- `group_replication_flow_control_release_percent`: Defines how the group quota should be released when flow control no longer needs to throttle the writer members. Added in MySQL 8.0.2.
- `Acl_cache_items_count`: Number of cached privilege objects. Added in MySQL 8.0.0.

- `binlog_expire_logs_seconds`: Purge binary logs after this many seconds; effects are additive with expire_logs_days. Added in MySQL 8.0.1.
- `binlog_row_metadata`: Configures the amount of table related metadata binary logged when using row-based logging.. Added in MySQL 8.0.1.
- `binlog_transaction_dependency_tracking`: Source of dependency information (commit timestamps or transaction write sets) from which to assess which transactions can be executed in parallel by slave's multi-threaded applier.. Added in MySQL 8.0.1.
- `binlog_transaction_dependency_history_size`: Number of row hashes kept for looking up transaction that last updated some row.. Added in MySQL 8.0.1.
- `Com_alter_user_default_role`: Count of ALTER USER ... DEFAULT ROLE statements. Added in MySQL 8.0.0.
- `Com_create_role`: Count of CREATE ROLE statements. Added in MySQL 8.0.0.
- `Com_drop_role`: Count of DROP ROLE statements. Added in MySQL 8.0.0.
- `Com_grant_roles`: Count of GRANT ROLE statements. Added in MySQL 8.0.0.
- `Com_install_component`: Count of INSTALL COMPONENT statements. Added in MySQL 8.0.0.
- `Com_revoke_roles`: Count of REVOKE ROLES statements. Added in MySQL 8.0.0.
- `Com_set_role`: Count of SET ROLE statements. Added in MySQL 8.0.0.
- `Com_uninstall_component`: Count of UINSTALL COMPONENT statements. Added in MySQL 8.0.0.
- `histogram_generation_max_mem_size`: Maximum memory for creating histogram statistics. Added in MySQL 8.0.2.
- `information_schema_stats`: Which table statistics source to use. Added in MySQL 8.0.0.
- `innodb_buffer_pool_debug`: Permits multiple buffer pool instances when the buffer pool is less than 1GB in size. Added in MySQL 8.0.0.
- `innodb_checkpoint_disabled`: Disables checkpoints so that a deliberate server exit always initiates recovery. Added in MySQL 8.0.2.
- `innodb_redo_log_encrypt`: Controls encryption of redo log data for encrypted tablepaces. Added in MySQL 8.0.1.
- `innodb_scan_directories`: Defines directories to scan for tablespace files during InnoDB recovery. Added in MySQL 8.0.2.
- `internal_tmp_mem_storage_engine`: Defines the storage to use for internal in-memory temporary tables. Added in MySQL 8.0.2.
- `innodb_undo_log_encrypt`: Controls encryption of undo log data for encrypted tablepaces. Added in MySQL 8.0.1.
- `original_commit_timestamp`: The time when a transaction was committed on the original master. Added in MySQL 8.0.1.
- `performance_schema_error_size`: Number of instrumented errors. Added in MySQL 8.0.0.
- `Performance_schema_session_connect_attrs_longest_seen`: The size of the longest valid connection attribute buffer seen. Added in MySQL 8.0.0.
- `persisted_globals_load`: Whether to load persisted configuration settings. Added in MySQL 8.0.0.

- `schema_definition_cache`: Defines the number of schema definition objects that can be kept in the dictionary object cache. Added in MySQL 8.0.0.
- `stored_program_definition_cache`: Defines the number of stored program and event definition objects that can be kept in the dictionary object cache. Added in MySQL 8.0.0.
- `tablespace_definition_cache`: Defines the number of tablespace definition objects that can be kept in the dictionary object cache. Added in MySQL 8.0.0.
- `temptable_max_ram`: Defines the maximum amount of memory that can occupied by the TempTable storage engine before data is stored on disk. Added in MySQL 8.0.2.

Options and Variables Deprecated in MySQL 8.0

The following system variables, status variables, and options have been deprecated in MySQL 8.0.

- `symbolic-links`: Permit symbolic links for MyISAM tables. Deprecated as of MySQL 8.0.2.

Options and Variables Removed in MySQL 8.0

The following system variables, status variables, and options have been removed in MySQL 8.0.

- `bootstrap`: Used by mysql installation scripts. Removed in MySQL 8.0.0.
- `Com_alter_db_upgrade`: Count of ALTER DATABASE ... UPGRADE DATA DIRECTORY NAME statements. Removed in MySQL 8.0.0.
- `ignore-db-dir`: Treat directory as nondatabase directory. Removed in MySQL 8.0.0.
- `ignore_db_dirs`: Directories treated as nondatabase directories. Removed in MySQL 8.0.0.
- `Innodb_available_undo_logs`: Display the total number of InnoDB rollback segments; different from innodb_rollback_segments, which displays the number of active rollback segments. Removed in MySQL 8.0.2.
- `innodb_checksums`: Enable InnoDB checksums validation. Removed in MySQL 8.0.0.
- `innodb_disable_resize_buffer_pool_debug`: Disables resizing of the InnoDB buffer pool. Removed in MySQL 8.0.0.
- `innodb_file_format`: The format for new InnoDB tables. Removed in MySQL 8.0.0.
- `innodb_file_format_check`: Whether InnoDB performs file format compatibility checking. Removed in MySQL 8.0.0.
- `innodb_file_format_max`: The file format tag in the shared tablespace. Removed in MySQL 8.0.0.
- `innodb_large_prefix`: Enables longer keys for column prefix indexes. Removed in MySQL 8.0.0.
- `innodb_locks_unsafe_for_binlog`: Force InnoDB not to use next-key locking. Instead use only row-level locking. Removed in MySQL 8.0.0.
- `innodb_stats_sample_pages`: Number of index pages to sample for index distribution statistics. Removed in MySQL 8.0.0.
- `innodb_support_xa`: Enable InnoDB support for the XA two-phase commit. Removed in MySQL 8.0.0.
- `innodb_undo_logs`: Defines the number of undo logs (rollback segments) used by InnoDB; an alias for innodb_rollback_segments. Removed in MySQL 8.0.2.
- `partition`: Enable (or disable) partitioning support. Removed in MySQL 8.0.0.

- `show_compatibility_56`: Compatibility for SHOW STATUS/VARIABLES. Removed in MySQL 8.0.1.
- `skip-partition`: Do not enable user-defined partitioning. Removed in MySQL 8.0.0.
- `Slave_heartbeat_period`: The slave's replication heartbeat interval, in seconds. Removed in MySQL 8.0.1.
- `Slave_last_heartbeat`: Shows when the latest heartbeat signal was received, in TIMESTAMP format. Removed in MySQL 8.0.1.
- `Slave_received_heartbeats`: Number of heartbeats received by a replication slave since previous reset. Removed in MySQL 8.0.1.
- `Slave_retried_transactions`: The total number of times since startup that the replication slave SQL thread has retried transactions. Removed in MySQL 8.0.1.
- `Slave_running`: The state of this server as a replication slave (slave I/O thread status). Removed in MySQL 8.0.1.
- `sync_frm`: Sync .frm to disk on create. Enabled by default. Removed in MySQL 8.0.0.
- `temp-pool`: Using this option will cause most temporary files created to use a small set of names, rather than a unique name for each new file. Removed in MySQL 8.0.1.

1.6 MySQL Information Sources

This section lists sources of additional information that you may find helpful, such as MySQL web sites, mailing lists, user forums, and Internet Relay Chat.

1.6.1 MySQL Web Sites

The primary web site for MySQL documentation is http://dev.mysql.com/doc/. Online and downloadable documentation formats are available for the MySQL Reference Manual, MySQL Connectors, and more.

The MySQL developers provide information about new and upcoming features as the MySQL Server Blog.

1.6.2 MySQL Mailing Lists

This section introduces the MySQL mailing lists and provides guidelines as to how the lists should be used. When you subscribe to a mailing list, you receive all postings to the list as email messages. You can also send your own questions and answers to the list.

To subscribe to or unsubscribe from any of the mailing lists described in this section, visit http://lists.mysql.com/. For most of them, you can select the regular version of the list where you get individual messages, or a digest version where you get one large message per day.

Please *do not* send messages about subscribing or unsubscribing to any of the mailing lists, because such messages are distributed automatically to thousands of other users.

Your local site may have many subscribers to a MySQL mailing list. If so, the site may have a local mailing list, so that messages sent from `lists.mysql.com` to your site are propagated to the local list. In such cases, please contact your system administrator to be added to or dropped from the local MySQL list.

To have traffic for a mailing list go to a separate mailbox in your mail program, set up a filter based on the message headers. You can use either the `List-ID:` or `Delivered-To:` headers to identify list messages.

The MySQL mailing lists are as follows:

- `announce`

The list for announcements of new versions of MySQL and related programs. This is a low-volume list to which all MySQL users should subscribe.

- `mysql`

  The main list for general MySQL discussion. Please note that some topics are better discussed on the more-specialized lists. If you post to the wrong list, you may not get an answer.

- `bugs`

  The list for people who want to stay informed about issues reported since the last release of MySQL or who want to be actively involved in the process of bug hunting and fixing. See Section 1.7, “How to Report Bugs or Problems”.

- `internals`

  The list for people who work on the MySQL code. This is also the forum for discussions on MySQL development and for posting patches.

- `mysqldoc`

  The list for people who work on the MySQL documentation.

- `benchmarks`

  The list for anyone interested in performance issues. Discussions concentrate on database performance (not limited to MySQL), but also include broader categories such as performance of the kernel, file system, disk system, and so on.

- `packagers`

  The list for discussions on packaging and distributing MySQL. This is the forum used by distribution maintainers to exchange ideas on packaging MySQL and on ensuring that MySQL looks and feels as similar as possible on all supported platforms and operating systems.

- `java`

  The list for discussions about the MySQL server and Java. It is mostly used to discuss JDBC drivers such as MySQL Connector/J.

- `win32`

  The list for all topics concerning the MySQL software on Microsoft operating systems, such as Windows 9x, Me, NT, 2000, XP, and 2003.

- `myodbc`

  The list for all topics concerning connecting to the MySQL server with ODBC.

- `gui-tools`

  The list for all topics concerning MySQL graphical user interface tools such as MySQL Workbench.

- `cluster`

  The list for discussion of MySQL Cluster.

- `dotnet`

  The list for discussion of the MySQL server and the .NET platform. It is mostly related to MySQL Connector/Net.

- `plusplus`

The list for all topics concerning programming with the C++ API for MySQL.

- `perl`

  The list for all topics concerning Perl support for MySQL with `DBD::mysql`.

If you're unable to get an answer to your questions from a MySQL mailing list or forum, one option is to purchase support from Oracle. This puts you in direct contact with MySQL developers.

The following MySQL mailing lists are in languages other than English. These lists are not operated by Oracle.

- `<mysql-france-subscribe@yahoogroups.com>`

  A French mailing list.

- `<list@tinc.net>`

  A Korean mailing list. To subscribe, email `subscribe mysql your@email.address` to this list.

- `<mysql-de-request@lists.4t2.com>`

  A German mailing list. To subscribe, email `subscribe mysql-de your@email.address` to this list. You can find information about this mailing list at http://www.4t2.com/mysql/.

- `<mysql-br-request@listas.linkway.com.br>`

  A Portuguese mailing list. To subscribe, email `subscribe mysql-br your@email.address` to this list.

- `<mysql-alta@elistas.net>`

  A Spanish mailing list. To subscribe, email `subscribe mysql your@email.address` to this list.

1.6.2.1 Guidelines for Using the Mailing Lists

Please do not post mail messages from your browser with HTML mode turned on. Many users do not read mail with a browser.

When you answer a question sent to a mailing list, if you consider your answer to have broad interest, you may want to post it to the list instead of replying directly to the individual who asked. Try to make your answer general enough that people other than the original poster may benefit from it. When you post to the list, please make sure that your answer is not a duplication of a previous answer.

Try to summarize the essential part of the question in your reply. Do not feel obliged to quote the entire original message.

When answers are sent to you individually and not to the mailing list, it is considered good etiquette to summarize the answers and send the summary to the mailing list so that others may have the benefit of responses you received that helped you solve your problem.

1.6.3 MySQL Community Support at the MySQL Forums

The forums at http://forums.mysql.com are an important community resource. Many forums are available, grouped into these general categories:

- Migration
- MySQL Usage
- MySQL Connectors
- Programming Languages

- Tools
- 3rd-Party Applications
- Storage Engines
- MySQL Technology
- SQL Standards
- Business

1.6.4 MySQL Community Support on Internet Relay Chat (IRC)

In addition to the various MySQL mailing lists and forums, you can find experienced community people on Internet Relay Chat (IRC). These are the best networks/channels currently known to us:

freenode (see http://www.freenode.net/ for servers)

- `#mysql` is primarily for MySQL questions, but other database and general SQL questions are welcome. Questions about PHP, Perl, or C in combination with MySQL are also common.
- `#workbench` is primarily for MySQL Workbench related questions and thoughts, and it is also a good place to meet the MySQL Workbench developers.

If you are looking for IRC client software to connect to an IRC network, take a look at `xChat` (http://www.xchat.org/). X-Chat (GPL licensed) is available for Unix as well as for Windows platforms (a free Windows build of X-Chat is available at http://www.silverex.org/download/).

1.6.5 MySQL Enterprise

Oracle offers technical support in the form of MySQL Enterprise. For organizations that rely on the MySQL DBMS for business-critical production applications, MySQL Enterprise is a commercial subscription offering which includes:

- MySQL Enterprise Server
- MySQL Enterprise Monitor
- Monthly Rapid Updates and Quarterly Service Packs
- MySQL Knowledge Base
- 24x7 Technical and Consultative Support

MySQL Enterprise is available in multiple tiers, giving you the flexibility to choose the level of service that best matches your needs. For more information, see MySQL Enterprise.

1.7 How to Report Bugs or Problems

Before posting a bug report about a problem, please try to verify that it is a bug and that it has not been reported already:

- Start by searching the MySQL online manual at http://dev.mysql.com/doc/. We try to keep the manual up to date by updating it frequently with solutions to newly found problems. In addition, the release notes accompanying the manual can be particularly useful since it is quite possible that a newer version contains a solution to your problem. The release notes are available at the location just given for the manual.
- If you get a parse error for an SQL statement, please check your syntax closely. If you cannot find something wrong with it, it is extremely likely that your current version of MySQL Server doesn't support the syntax you are using. If you are using the current version and the manual doesn't cover the syntax that you are using, MySQL Server doesn't support your statement.

If the manual covers the syntax you are using, but you have an older version of MySQL Server, you should check the MySQL change history to see when the syntax was implemented. In this case, you have the option of upgrading to a newer version of MySQL Server.

- For solutions to some common problems, see Section B.5, “Problems and Common Errors”.
- Search the bugs database at http://bugs.mysql.com/ to see whether the bug has been reported and fixed.
- Search the MySQL mailing list archives at http://lists.mysql.com/. See Section 1.6.2, “MySQL Mailing Lists”.
- You can also use http://www.mysql.com/search/ to search all the Web pages (including the manual) that are located at the MySQL Web site.

If you cannot find an answer in the manual, the bugs database, or the mailing list archives, check with your local MySQL expert. If you still cannot find an answer to your question, please use the following guidelines for reporting the bug.

The normal way to report bugs is to visit http://bugs.mysql.com/, which is the address for our bugs database. This database is public and can be browsed and searched by anyone. If you log in to the system, you can enter new reports.

Bugs posted in the bugs database at http://bugs.mysql.com/ that are corrected for a given release are noted in the release notes.

If you find a sensitive security bug in MySQL Server, please let us know immediately by sending an email message to `<secalert_us@oracle.com>`. Exception: Support customers should report all problems, including security bugs, to Oracle Support at http://support.oracle.com/.

To discuss problems with other users, you can use one of the MySQL mailing lists. Section 1.6.2, “MySQL Mailing Lists”.

Writing a good bug report takes patience, but doing it right the first time saves time both for us and for yourself. A good bug report, containing a full test case for the bug, makes it very likely that we will fix the bug in the next release. This section helps you write your report correctly so that you do not waste your time doing things that may not help us much or at all. Please read this section carefully and make sure that all the information described here is included in your report.

Preferably, you should test the problem using the latest production or development version of MySQL Server before posting. Anyone should be able to repeat the bug by just using `mysql test < script_file` on your test case or by running the shell or Perl script that you include in the bug report. Any bug that we are able to repeat has a high chance of being fixed in the next MySQL release.

It is most helpful when a good description of the problem is included in the bug report. That is, give a good example of everything you did that led to the problem and describe, in exact detail, the problem itself. The best reports are those that include a full example showing how to reproduce the bug or problem. See Section 28.5, “Debugging and Porting MySQL”.

Remember that it is possible for us to respond to a report containing too much information, but not to one containing too little. People often omit facts because they think they know the cause of a problem and assume that some details do not matter. A good principle to follow is that if you are in doubt about stating something, state it. It is faster and less troublesome to write a couple more lines in your report than to wait longer for the answer if we must ask you to provide information that was missing from the initial report.

The most common errors made in bug reports are (a) not including the version number of the MySQL distribution that you use, and (b) not fully describing the platform on which the MySQL server is installed (including the platform type and version number). These are highly relevant pieces of information, and in 99 cases out of 100, the bug report is useless without them. Very often we get questions like, “Why doesn't this work for me?” Then we find that the feature requested wasn't

implemented in that MySQL version, or that a bug described in a report has been fixed in newer MySQL versions. Errors often are platform-dependent. In such cases, it is next to impossible for us to fix anything without knowing the operating system and the version number of the platform.

If you compiled MySQL from source, remember also to provide information about your compiler if it is related to the problem. Often people find bugs in compilers and think the problem is MySQL-related. Most compilers are under development all the time and become better version by version. To determine whether your problem depends on your compiler, we need to know what compiler you used. Note that every compiling problem should be regarded as a bug and reported accordingly.

If a program produces an error message, it is very important to include the message in your report. If we try to search for something from the archives, it is better that the error message reported exactly matches the one that the program produces. (Even the lettercase should be observed.) It is best to copy and paste the entire error message into your report. You should never try to reproduce the message from memory.

If you have a problem with Connector/ODBC (MyODBC), please try to generate a trace file and send it with your report. See How to Report Connector/ODBC Problems or Bugs.

If your report includes long query output lines from test cases that you run with the `mysql` command-line tool, you can make the output more readable by using the `--vertical` option or the `\G` statement terminator. The `EXPLAIN SELECT` example later in this section demonstrates the use of `\G`.

Please include the following information in your report:

- The version number of the MySQL distribution you are using (for example, MySQL 5.7.10). You can find out which version you are running by executing `mysqladmin version`. The `mysqladmin` program can be found in the `bin` directory under your MySQL installation directory.
- The manufacturer and model of the machine on which you experience the problem.
- The operating system name and version. If you work with Windows, you can usually get the name and version number by double-clicking your My Computer icon and pulling down the “Help/About Windows” menu. For most Unix-like operating systems, you can get this information by executing the command `uname -a`.
- Sometimes the amount of memory (real and virtual) is relevant. If in doubt, include these values.
- If you are using a source distribution of the MySQL software, include the name and version number of the compiler that you used. If you have a binary distribution, include the distribution name.
- If the problem occurs during compilation, include the exact error messages and also a few lines of context around the offending code in the file where the error occurs.
- If `mysqld` died, you should also report the statement that crashed `mysqld`. You can usually get this information by running `mysqld` with query logging enabled, and then looking in the log after `mysqld` crashes. See Section 28.5, “Debugging and Porting MySQL”.
- If a database table is related to the problem, include the output from the `SHOW CREATE TABLE` *`db_name.tbl_name`* statement in the bug report. This is a very easy way to get the definition of any table in a database. The information helps us create a situation matching the one that you have experienced.
- The SQL mode in effect when the problem occurred can be significant, so please report the value of the `sql_mode` system variable. For stored procedure, stored function, and trigger objects, the relevant `sql_mode` value is the one in effect when the object was created. For a stored procedure or function, the `SHOW CREATE PROCEDURE` or `SHOW CREATE FUNCTION` statement shows the relevant SQL mode, or you can query `INFORMATION_SCHEMA` for the information:

```
SELECT ROUTINE_SCHEMA, ROUTINE_NAME, SQL_MODE
FROM INFORMATION_SCHEMA.ROUTINES;
```

For triggers, you can use this statement:

```
SELECT EVENT_OBJECT_SCHEMA, EVENT_OBJECT_TABLE, TRIGGER_NAME, SQL_MODE
FROM INFORMATION_SCHEMA.TRIGGERS;
```

- For performance-related bugs or problems with `SELECT` statements, you should always include the output of `EXPLAIN SELECT ...`, and at least the number of rows that the `SELECT` statement produces. You should also include the output from `SHOW CREATE TABLE tbl_name` for each table that is involved. The more information you provide about your situation, the more likely it is that someone can help you.

  The following is an example of a very good bug report. The statements are run using the `mysql` command-line tool. Note the use of the `\G` statement terminator for statements that would otherwise provide very long output lines that are difficult to read.

```
mysql> SHOW VARIABLES;
mysql> SHOW COLUMNS FROM ...\G
       <output from SHOW COLUMNS>
mysql> EXPLAIN SELECT ...\G
       <output from EXPLAIN>
mysql> FLUSH STATUS;
mysql> SELECT ...;
       <A short version of the output from SELECT,
       including the time taken to run the query>
mysql> SHOW STATUS;
       <output from SHOW STATUS>
```

- If a bug or problem occurs while running `mysqld`, try to provide an input script that reproduces the anomaly. This script should include any necessary source files. The more closely the script can reproduce your situation, the better. If you can make a reproducible test case, you should upload it to be attached to the bug report.

  If you cannot provide a script, you should at least include the output from `mysqladmin variables extended-status processlist` in your report to provide some information on how your system is performing.

- If you cannot produce a test case with only a few rows, or if the test table is too big to be included in the bug report (more than 10 rows), you should dump your tables using `mysqldump` and create a `README` file that describes your problem. Create a compressed archive of your files using `tar` and `gzip` or `zip`. After you initiate a bug report for our bugs database at http://bugs.mysql.com/, click the Files tab in the bug report for instructions on uploading the archive to the bugs database.

- If you believe that the MySQL server produces a strange result from a statement, include not only the result, but also your opinion of what the result should be, and an explanation describing the basis for your opinion.

- When you provide an example of the problem, it is better to use the table names, variable names, and so forth that exist in your actual situation than to come up with new names. The problem could be related to the name of a table or variable. These cases are rare, perhaps, but it is better to be safe than sorry. After all, it should be easier for you to provide an example that uses your actual situation, and it is by all means better for us. If you have data that you do not want to be visible to others in the bug report, you can upload it using the Files tab as previously described. If the information is really top secret and you do not want to show it even to us, go ahead and provide an example using other names, but please regard this as the last choice.

- Include all the options given to the relevant programs, if possible. For example, indicate the options that you use when you start the `mysqld` server, as well as the options that you use to run any MySQL client programs. The options to programs such as `mysqld` and `mysql`, and to the `configure` script, are often key to resolving problems and are very relevant. It is never a bad idea to include them. If your problem involves a program written in a language such as Perl or PHP,

please include the language processor's version number, as well as the version for any modules that the program uses. For example, if you have a Perl script that uses the `DBI` and `DBD::mysql` modules, include the version numbers for Perl, `DBI`, and `DBD::mysql`.

- If your question is related to the privilege system, please include the output of `mysqladmin reload`, and all the error messages you get when trying to connect. When you test your privileges, you should execute `mysqladmin reload version` and try to connect with the program that gives you trouble.
- If you have a patch for a bug, do include it. But do not assume that the patch is all we need, or that we can use it, if you do not provide some necessary information such as test cases showing the bug that your patch fixes. We might find problems with your patch or we might not understand it at all. If so, we cannot use it.

  If we cannot verify the exact purpose of the patch, we will not use it. Test cases help us here. Show that the patch handles all the situations that may occur. If we find a borderline case (even a rare one) where the patch will not work, it may be useless.
- Guesses about what the bug is, why it occurs, or what it depends on are usually wrong. Even the MySQL team cannot guess such things without first using a debugger to determine the real cause of a bug.
- Indicate in your bug report that you have checked the reference manual and mail archive so that others know you have tried to solve the problem yourself.
- If your data appears corrupt or you get errors when you access a particular table, first check your tables with `CHECK TABLE`. If that statement reports any errors:
  - The `InnoDB` crash recovery mechanism handles cleanup when the server is restarted after being killed, so in typical operation there is no need to “repair” tables. If you encounter an error with `InnoDB` tables, restart the server and see whether the problem persists, or whether the error affected only cached data in memory. If data is corrupted on disk, consider restarting with the `innodb_force_recovery` option enabled so that you can dump the affected tables.
  - For non-transactional tables, try to repair them with `REPAIR TABLE` or with `myisamchk`. See Chapter 5, *MySQL Server Administration*.

  If you are running Windows, please verify the value of `lower_case_table_names` using the `SHOW VARIABLES LIKE 'lower_case_table_names'` statement. This variable affects how the server handles lettercase of database and table names. Its effect for a given value should be as described in Section 9.2.2, “Identifier Case Sensitivity”.
- If you often get corrupted tables, you should try to find out when and why this happens. In this case, the error log in the MySQL data directory may contain some information about what happened. (This is the file with the `.err` suffix in the name.) See Section 5.4.2, “The Error Log”. Please include any relevant information from this file in your bug report. Normally `mysqld` should *never* crash a table if nothing killed it in the middle of an update. If you can find the cause of `mysqld` dying, it is much easier for us to provide you with a fix for the problem. See Section B.5.1, “How to Determine What Is Causing a Problem”.
- If possible, download and install the most recent version of MySQL Server and check whether it solves your problem. All versions of the MySQL software are thoroughly tested and should work without problems. We believe in making everything as backward-compatible as possible, and you should be able to switch MySQL versions without difficulty. See Section 2.1.1, “Which MySQL Version and Distribution to Install”.

1.8 MySQL Standards Compliance

This section describes how MySQL relates to the ANSI/ISO SQL standards. MySQL Server has many extensions to the SQL standard, and here you can find out what they are and how to use them. You

can also find information about functionality missing from MySQL Server, and how to work around some of the differences.

The SQL standard has been evolving since 1986 and several versions exist. In this manual, “SQL-92” refers to the standard released in 1992, “SQL:1999” refers to the standard released in 1999, “SQL:2003” refers to the standard released in 2003, and “SQL:2008” refers to the most recent version of the standard, released in 2008. We use the phrase “the SQL standard” or “standard SQL” to mean the current version of the SQL Standard at any time.

One of our main goals with the product is to continue to work toward compliance with the SQL standard, but without sacrificing speed or reliability. We are not afraid to add extensions to SQL or support for non-SQL features if this greatly increases the usability of MySQL Server for a large segment of our user base. The `HANDLER` interface is an example of this strategy. See Section 13.2.4, “HANDLER Syntax”.

We continue to support transactional and nontransactional databases to satisfy both mission-critical 24/7 usage and heavy Web or logging usage.

MySQL Server was originally designed to work with medium-sized databases (10-100 million rows, or about 100MB per table) on small computer systems. Today MySQL Server handles terabyte-sized databases.

We are not targeting real-time support, although MySQL replication capabilities offer significant functionality.

MySQL supports ODBC levels 0 to 3.51.

MySQL supports high-availability database clustering using the `NDBCLUSTER` storage engine. See MySQL NDB Cluster 7.5 and NDB Cluster 7.6.

We implement XML functionality which supports most of the W3C XPath standard. See Section 12.11, “XML Functions”.

MySQL supports a native JSON data type as defined by RFC 7159, and based on the ECMAScript standard (ECMA-262). See Section 11.6, “The JSON Data Type”. MySQL also implements a subset of the SQL/JSON functions specified by a pre-publication draft of the SQL:2016 standard; see Section 12.16, “JSON Functions”, for more information.

Selecting SQL Modes

The MySQL server can operate in different SQL modes, and can apply these modes differently for different clients, depending on the value of the `sql_mode` system variable. DBAs can set the global SQL mode to match site server operating requirements, and each application can set its session SQL mode to its own requirements.

Modes affect the SQL syntax MySQL supports and the data validation checks it performs. This makes it easier to use MySQL in different environments and to use MySQL together with other database servers.

For more information on setting the SQL mode, see Section 5.1.8, “Server SQL Modes”.

Running MySQL in ANSI Mode

To run MySQL Server in ANSI mode, start `mysqld` with the `--ansi` option. Running the server in ANSI mode is the same as starting it with the following options:

```
--transaction-isolation=SERIALIZABLE --sql-mode=ANSI
```

To achieve the same effect at runtime, execute these two statements:

```
SET GLOBAL TRANSACTION ISOLATION LEVEL SERIALIZABLE;
SET GLOBAL sql_mode = 'ANSI';
```

You can see that setting the `sql_mode` system variable to `'ANSI'` enables all SQL mode options that are relevant for ANSI mode as follows:

```
mysql> SET GLOBAL sql_mode='ANSI';
mysql> SELECT @@global.sql_mode;
       -> 'REAL_AS_FLOAT,PIPES_AS_CONCAT,ANSI_QUOTES,IGNORE_SPACE,ANSI'
```

Running the server in ANSI mode with `--ansi` is not quite the same as setting the SQL mode to `'ANSI'` because the `--ansi` option also sets the transaction isolation level.

See Section 5.1.4, “Server Command Options”.

1.8.1 MySQL Extensions to Standard SQL

MySQL Server supports some extensions that you probably will not find in other SQL DBMSs. Be warned that if you use them, your code will not be portable to other SQL servers. In some cases, you can write code that includes MySQL extensions, but is still portable, by using comments of the following form:

```
/*! MySQL-specific code */
```

In this case, MySQL Server parses and executes the code within the comment as it would any other SQL statement, but other SQL servers will ignore the extensions. For example, MySQL Server recognizes the `STRAIGHT_JOIN` keyword in the following statement, but other servers will not:

```
SELECT /*! STRAIGHT_JOIN */ col1 FROM table1,table2 WHERE ...
```

If you add a version number after the `!` character, the syntax within the comment is executed only if the MySQL version is greater than or equal to the specified version number. The `KEY_BLOCK_SIZE` clause in the following comment is executed only by servers from MySQL 5.1.10 or higher:

```
CREATE TABLE t1(a INT, KEY (a)) /*!50110 KEY_BLOCK_SIZE=1024 */;
```

The following descriptions list MySQL extensions, organized by category.

- Organization of data on disk

  MySQL Server maps each database to a directory under the MySQL data directory, and maps tables within a database to file names in the database directory. Consequently, database and table names are case sensitive in MySQL Server on operating systems that have case-sensitive file names (such as most Unix systems). See Section 9.2.2, “Identifier Case Sensitivity”.

- General language syntax

  - By default, strings can be enclosed by `"` as well as `'`. If the `ANSI_QUOTES` SQL mode is enabled, strings can be enclosed only by `'` and the server interprets strings enclosed by `"` as identifiers.
  - `\` is the escape character in strings.
  - In SQL statements, you can access tables from different databases with the *`db_name.tbl_name`* syntax. Some SQL servers provide the same functionality but call this `User space`. MySQL Server doesn't support tablespaces such as used in statements like this: `CREATE TABLE ralph.my_table ... IN my_tablespace`.

- SQL statement syntax

  - The `ANALYZE TABLE`, `CHECK TABLE`, `OPTIMIZE TABLE`, and `REPAIR TABLE` statements.
  - The `CREATE DATABASE`, `DROP DATABASE`, and `ALTER DATABASE` statements. See Section 13.1.9, “CREATE DATABASE Syntax”, Section 13.1.19, “DROP DATABASE Syntax”, and Section 13.1.1, “ALTER DATABASE Syntax”.

- The `DO` statement.
- `EXPLAIN SELECT` to obtain a description of how tables are processed by the query optimizer.
- The `FLUSH` and `RESET` statements.
- The `SET` statement. See Section 13.7.4.1, “SET Syntax for Variable Assignment”.
- The `SHOW` statement. See Section 13.7.5, “SHOW Syntax”. The information produced by many of the MySQL-specific `SHOW` statements can be obtained in more standard fashion by using `SELECT` to query `INFORMATION_SCHEMA`. See Chapter 24, *INFORMATION_SCHEMA Tables*.
- Use of `LOAD DATA INFILE`. In many cases, this syntax is compatible with Oracle's `LOAD DATA INFILE`. See Section 13.2.7, “LOAD DATA INFILE Syntax”.
- Use of `RENAME TABLE`. See Section 13.1.29, “RENAME TABLE Syntax”.
- Use of `REPLACE` instead of `DELETE` plus `INSERT`. See Section 13.2.9, “REPLACE Syntax”.
- Use of `CHANGE col_name`, `DROP col_name`, or `DROP INDEX`, `IGNORE` or `RENAME` in `ALTER TABLE` statements. Use of multiple `ADD`, `ALTER`, `DROP`, or `CHANGE` clauses in an `ALTER TABLE` statement. See Section 13.1.7, “ALTER TABLE Syntax”.
- Use of index names, indexes on a prefix of a column, and use of `INDEX` or `KEY` in `CREATE TABLE` statements. See Section 13.1.15, “CREATE TABLE Syntax”.
- Use of `TEMPORARY` or `IF NOT EXISTS` with `CREATE TABLE`.
- Use of `IF EXISTS` with `DROP TABLE` and `DROP DATABASE`.
- The capability of dropping multiple tables with a single `DROP TABLE` statement.
- The `ORDER BY` and `LIMIT` clauses of the `UPDATE` and `DELETE` statements.
- `INSERT INTO tbl_name SET col_name = ...` syntax.
- The `DELAYED` clause of the `INSERT` and `REPLACE` statements.
- The `LOW_PRIORITY` clause of the `INSERT`, `REPLACE`, `DELETE`, and `UPDATE` statements.
- Use of `INTO OUTFILE` or `INTO DUMPFILE` in `SELECT` statements. See Section 13.2.10, “SELECT Syntax”.
- Options such as `STRAIGHT_JOIN` or `SQL_SMALL_RESULT` in `SELECT` statements.
- You don't need to name all selected columns in the `GROUP BY` clause. This gives better performance for some very specific, but quite normal queries. See Section 12.18, “Aggregate (GROUP BY) Functions”.
- You can specify `ASC` and `DESC` with `GROUP BY`, not just with `ORDER BY`.
- The ability to set variables in a statement with the `:=` assignment operator. See Section 9.4, “User-Defined Variables”.

- Data types
  - The `MEDIUMINT`, `SET`, and `ENUM` data types, and the various `BLOB` and `TEXT` data types.
  - The `AUTO_INCREMENT`, `BINARY`, `NULL`, `UNSIGNED`, and `ZEROFILL` data type attributes.
- Functions and operators

- To make it easier for users who migrate from other SQL environments, MySQL Server supports aliases for many functions. For example, all string functions support both standard SQL syntax and ODBC syntax.
- MySQL Server understands the `||` and `&&` operators to mean logical OR and AND, as in the C programming language. In MySQL Server, `||` and `OR` are synonyms, as are `&&` and `AND`. Because of this nice syntax, MySQL Server doesn't support the standard SQL `||` operator for string concatenation; use `CONCAT()` instead. Because `CONCAT()` takes any number of arguments, it is easy to convert use of the `||` operator to MySQL Server.
- Use of `COUNT(DISTINCT value_list)` where *`value_list`* has more than one element.
- String comparisons are case insensitive by default, with sort ordering determined by the collation of the current character set, which is `utf8mb4` by default. To perform case-sensitive comparisons instead, you should declare your columns with the `BINARY` attribute or use the `BINARY` cast, which causes comparisons to be done using the underlying character code values rather than a lexical ordering.
- The `%` operator is a synonym for `MOD()`. That is, *`N % M`* is equivalent to `MOD(N,M)`. `%` is supported for C programmers and for compatibility with PostgreSQL.
- The `=`, `<>`, `<=`, `<`, `>=`, `>`, `<<`, `>>`, `<=>`, `AND`, `OR`, or `LIKE` operators may be used in expressions in the output column list (to the left of the `FROM`) in `SELECT` statements. For example:

```
mysql> SELECT col1=1 AND col2=2 FROM my_table;
```

- The `LAST_INSERT_ID()` function returns the most recent `AUTO_INCREMENT` value. See Section 12.14, “Information Functions”.
- `LIKE` is permitted on numeric values.
- The `REGEXP` and `NOT REGEXP` extended regular expression operators.
- `CONCAT()` or `CHAR()` with one argument or more than two arguments. (In MySQL Server, these functions can take a variable number of arguments.)
- The `BIT_COUNT()`, `CASE`, `ELT()`, `FROM_DAYS()`, `FORMAT()`, `IF()`, `PASSWORD()`, `ENCRYPT()`, `MD5()`, `ENCODE()`, `DECODE()`, `PERIOD_ADD()`, `PERIOD_DIFF()`, `TO_DAYS()`, and `WEEKDAY()` functions.
- Use of `TRIM()` to trim substrings. Standard SQL supports removal of single characters only.
- The `GROUP BY` functions `STD()`, `BIT_OR()`, `BIT_AND()`, `BIT_XOR()`, and `GROUP_CONCAT()`. See Section 12.18, “Aggregate (GROUP BY) Functions”.

1.8.2 MySQL Differences from Standard SQL

We try to make MySQL Server follow the ANSI SQL standard and the ODBC SQL standard, but MySQL Server performs operations differently in some cases:

- There are several differences between the MySQL and standard SQL privilege systems. For example, in MySQL, privileges for a table are not automatically revoked when you delete a table. You must explicitly issue a `REVOKE` statement to revoke privileges for a table. For more information, see Section 13.7.1.8, “REVOKE Syntax”.
- The `CAST()` function does not support cast to `REAL` or `BIGINT`. See Section 12.10, “Cast Functions and Operators”.

1.8.2.1 SELECT INTO TABLE Differences

MySQL Server doesn't support the `SELECT ... INTO TABLE` Sybase SQL extension. Instead, MySQL Server supports the `INSERT INTO ... SELECT` standard SQL syntax, which is basically the same thing. See Section 13.2.6.1, “INSERT ... SELECT Syntax”. For example:

```
INSERT INTO tbl_temp2 (fld_id)
    SELECT tbl_temp1.fld_order_id
    FROM tbl_temp1 WHERE tbl_temp1.fld_order_id > 100;
```

Alternatively, you can use `SELECT ... INTO OUTFILE` or `CREATE TABLE ... SELECT`.

You can use `SELECT ... INTO` with user-defined variables. The same syntax can also be used inside stored routines using cursors and local variables. See Section 13.2.10.1, “SELECT ... INTO Syntax”.

1.8.2.2 UPDATE Differences

If you access a column from the table to be updated in an expression, `UPDATE` uses the current value of the column. The second assignment in the following statement sets `col2` to the current (updated) `col1` value, not the original `col1` value. The result is that `col1` and `col2` have the same value. This behavior differs from standard SQL.

```
UPDATE t1 SET col1 = col1 + 1, col2 = col1;
```

1.8.2.3 Foreign Key Differences

The MySQL implementation of foreign keys differs from the SQL standard in the following key respects:

- If there are several rows in the parent table that have the same referenced key value, `InnoDB` acts in foreign key checks as if the other parent rows with the same key value do not exist. For example, if you have defined a `RESTRICT` type constraint, and there is a child row with several parent rows, `InnoDB` does not permit the deletion of any of those parent rows.

  `InnoDB` performs cascading operations through a depth-first algorithm, based on records in the indexes corresponding to the foreign key constraints.

- A `FOREIGN KEY` constraint that references a non-`UNIQUE` key is not standard SQL but rather an `InnoDB` extension.

- If `ON UPDATE CASCADE` or `ON UPDATE SET NULL` recurses to update the *same table* it has previously updated during the same cascade, it acts like `RESTRICT`. This means that you cannot use self-referential `ON UPDATE CASCADE` or `ON UPDATE SET NULL` operations. This is to prevent infinite loops resulting from cascaded updates. A self-referential `ON DELETE SET NULL`, on the other hand, is possible, as is a self-referential `ON DELETE CASCADE`. Cascading operations may not be nested more than 15 levels deep.

- In an SQL statement that inserts, deletes, or updates many rows, foreign key constraints (like unique constraints) are checked row-by-row. When performing foreign key checks, `InnoDB` sets shared row-level locks on child or parent records that it must examine. MySQL checks foreign key constraints immediately; the check is not deferred to transaction commit. According to the SQL standard, the default behavior should be deferred checking. That is, constraints are only checked after the *entire SQL statement* has been processed. This means that it is not possible to delete a row that refers to itself using a foreign key.

For information about how the `InnoDB` storage engine handles foreign keys, see Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

1.8.2.4 '--' as the Start of a Comment

Standard SQL uses the C syntax `/* this is a comment */` for comments, and MySQL Server supports this syntax as well. MySQL also support extensions to this syntax that enable MySQL-specific SQL to be embedded in the comment, as described in Section 9.6, “Comment Syntax”.

Standard SQL uses "`--`" as a start-comment sequence. MySQL Server uses `#` as the start comment character. MySQL Server also supports a variant of the `--` comment style. That is, the `--` start-comment sequence must be followed by a space (or by a control character such as a newline). The space is required to prevent problems with automatically generated SQL queries that use constructs such as the following, where we automatically insert the value of the payment for `payment`:

```
UPDATE account SET credit=credit-payment
```

Consider about what happens if `payment` has a negative value such as `-1`:

```
UPDATE account SET credit=credit--1
```

`credit--1` is a valid expression in SQL, but `--` is interpreted as the start of a comment, part of the expression is discarded. The result is a statement that has a completely different meaning than intended:

```
UPDATE account SET credit=credit
```

The statement produces no change in value at all. This illustrates that permitting comments to start with `--` can have serious consequences.

Using our implementation requires a space following the `--` for it to be recognized as a start-comment sequence in MySQL Server. Therefore, `credit--1` is safe to use.

Another safe feature is that the `mysql` command-line client ignores lines that start with `--`.

1.8.3 How MySQL Deals with Constraints

MySQL enables you to work both with transactional tables that permit rollback and with nontransactional tables that do not. Because of this, constraint handling is a bit different in MySQL than in other DBMSs. We must handle the case when you have inserted or updated a lot of rows in a nontransactional table for which changes cannot be rolled back when an error occurs.

The basic philosophy is that MySQL Server tries to produce an error for anything that it can detect while parsing a statement to be executed, and tries to recover from any errors that occur while executing the statement. We do this in most cases, but not yet for all.

The options MySQL has when an error occurs are to stop the statement in the middle or to recover as well as possible from the problem and continue. By default, the server follows the latter course. This means, for example, that the server may coerce invalid values to the closest valid values.

Several SQL mode options are available to provide greater control over handling of bad data values and whether to continue statement execution or abort when errors occur. Using these options, you can configure MySQL Server to act in a more traditional fashion that is like other DBMSs that reject improper input. The SQL mode can be set globally at server startup to affect all clients. Individual clients can set the SQL mode at runtime, which enables each client to select the behavior most appropriate for its requirements. See Section 5.1.8, "Server SQL Modes".

The following sections describe how MySQL Server handles different types of constraints.

1.8.3.1 PRIMARY KEY and UNIQUE Index Constraints

Normally, errors occur for data-change statements (such as `INSERT` or `UPDATE`) that would violate primary-key, unique-key, or foreign-key constraints. If you are using a transactional storage engine such as `InnoDB`, MySQL automatically rolls back the statement. If you are using a nontransactional storage engine, MySQL stops processing the statement at the row for which the error occurred and leaves any remaining rows unprocessed.

MySQL supports an `IGNORE` keyword for `INSERT`, `UPDATE`, and so forth. If you use it, MySQL ignores primary-key or unique-key violations and continues processing with the next row. See the section

for the statement that you are using (Section 13.2.6, “INSERT Syntax”, Section 13.2.12, “UPDATE Syntax”, and so forth).

You can get information about the number of rows actually inserted or updated with the `mysql_info()` C API function. You can also use the `SHOW WARNINGS` statement. See Section 27.7.7.36, “mysql_info()”, and Section 13.7.5.40, “SHOW WARNINGS Syntax”.

Only `InnoDB` tables support foreign keys. See Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

1.8.3.2 FOREIGN KEY Constraints

Foreign keys let you cross-reference related data across tables, and foreign key constraints help keep this spread-out data consistent.

MySQL supports `ON UPDATE` and `ON DELETE` foreign key references in `CREATE TABLE` and `ALTER TABLE` statements. The available referential actions are `RESTRICT` (the default), `CASCADE`, `SET NULL`, and `NO ACTION`.

`SET DEFAULT` is also supported by the MySQL Server but is currently rejected as invalid by `InnoDB`. Since MySQL does not support deferred constraint checking, `NO ACTION` is treated as `RESTRICT`. For the exact syntax supported by MySQL for foreign keys, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”.

`MATCH FULL`, `MATCH PARTIAL`, and `MATCH SIMPLE` are allowed, but their use should be avoided, as they cause the MySQL Server to ignore any `ON DELETE` or `ON UPDATE` clause used in the same statement. `MATCH` options do not have any other effect in MySQL, which in effect enforces `MATCH SIMPLE` semantics full-time.

MySQL requires that foreign key columns be indexed; if you create a table with a foreign key constraint but no index on a given column, an index is created.

You can obtain information about foreign keys from the `INFORMATION_SCHEMA.KEY_COLUMN_USAGE` table. An example of a query against this table is shown here:

```
mysql> SELECT TABLE_SCHEMA, TABLE_NAME, COLUMN_NAME, CONSTRAINT_NAME
     > FROM INFORMATION_SCHEMA.KEY_COLUMN_USAGE
     > WHERE REFERENCED_TABLE_SCHEMA IS NOT NULL;
+--------------+---------------+-------------+-----------------+
| TABLE_SCHEMA | TABLE_NAME    | COLUMN_NAME | CONSTRAINT_NAME |
+--------------+---------------+-------------+-----------------+
| fk1          | myuser        | myuser_id   | f               |
| fk1          | product_order | customer_id | f2              |
| fk1          | product_order | product_id  | f1              |
+--------------+---------------+-------------+-----------------+
3 rows in set (0.01 sec)
```

Information about foreign keys on `InnoDB` tables can also be found in the `INNODB_SYS_FOREIGN` and `INNODB_SYS_FOREIGN_COLS` tables, in the `INFORMATION_SCHEMA` database.

Only `InnoDB` tables support foreign keys. See Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”, for information specific to foreign key support in `InnoDB`.

1.8.3.3 Constraints on Invalid Data

By default, MySQL is forgiving of invalid or improper data values and coerces them to valid values for data entry. However, you can enable strict SQL mode to select more traditional treatment of bad values such that the server rejects them and aborts the statement in which they occur. See Section 5.1.8, “Server SQL Modes”.

This section describes the default (forgiving) behavior of MySQL, as well as the strict SQL mode and how it differs.

If you are not using strict mode, then whenever you insert an “incorrect” value into a column, such as a `NULL` into a `NOT NULL` column or a too-large numeric value into a numeric column, MySQL sets the column to the “best possible value” instead of producing an error: The following rules describe in more detail how this works:

- If you try to store an out of range value into a numeric column, MySQL Server instead stores zero, the smallest possible value, or the largest possible value, whichever is closest to the invalid value.
- For strings, MySQL stores either the empty string or as much of the string as can be stored in the column.
- If you try to store a string that does not start with a number into a numeric column, MySQL Server stores 0.
- Invalid values for `ENUM` and `SET` columns are handled as described in Section 1.8.3.4, “ENUM and SET Constraints”.
- MySQL permits you to store certain incorrect date values into `DATE` and `DATETIME` columns (such as `'2000-02-31'` or `'2000-02-00'`). In this case, when an application has not enabled strict SQL mode, it up to the application to validate the dates before storing them. If MySQL can store a date value and retrieve exactly the same value, MySQL stores it as given. If the date is totally wrong (outside the server's ability to store it), the special “zero” date value `'0000-00-00'` is stored in the column instead.
- If you try to store `NULL` into a column that doesn't take `NULL` values, an error occurs for single-row `INSERT` statements. For multiple-row `INSERT` statements or for `INSERT INTO ... SELECT` statements, MySQL Server stores the implicit default value for the column data type. In general, this is `0` for numeric types, the empty string (`''`) for string types, and the “zero” value for date and time types. Implicit default values are discussed in Section 11.7, “Data Type Default Values”.
- If an `INSERT` statement specifies no value for a column, MySQL inserts its default value if the column definition includes an explicit `DEFAULT` clause. If the definition has no such `DEFAULT` clause, MySQL inserts the implicit default value for the column data type.

The reason for using the preceding rules in nonstrict mode is that we can't check these conditions until the statement has begun executing. We can't just roll back if we encounter a problem after updating a few rows, because the storage engine may not support rollback. The option of terminating the statement is not that good; in this case, the update would be “half done,” which is probably the worst possible scenario. In this case, it is better to “do the best you can” and then continue as if nothing happened.

You can select stricter treatment of input values by using the `STRICT_TRANS_TABLES` or `STRICT_ALL_TABLES` SQL modes:

```
SET sql_mode = 'STRICT_TRANS_TABLES';
SET sql_mode = 'STRICT_ALL_TABLES';
```

`STRICT_TRANS_TABLES` enables strict mode for transactional storage engines, and also to some extent for nontransactional engines. It works like this:

- For transactional storage engines, bad data values occurring anywhere in a statement cause the statement to abort and roll back.
- For nontransactional storage engines, a statement aborts if the error occurs in the first row to be inserted or updated. (When the error occurs in the first row, the statement can be aborted to leave the table unchanged, just as for a transactional table.) Errors in rows after the first do not abort the statement, because the table has already been changed by the first row. Instead, bad data values are adjusted and result in warnings rather than errors. In other words, with `STRICT_TRANS_TABLES`, a wrong value causes MySQL to roll back all updates done so far, if that can be done without changing the table. But once the table has been changed, further errors result in adjustments and warnings.

For even stricter checking, enable `STRICT_ALL_TABLES`. This is the same as `STRICT_TRANS_TABLES` except that for nontransactional storage engines, errors abort the statement even for bad data in rows following the first row. This means that if an error occurs partway through a multiple-row insert or update for a nontransactional table, a partial update results. Earlier rows are inserted or updated, but those from the point of the error on are not. To avoid this for nontransactional tables, either use single-row statements or else use `STRICT_TRANS_TABLES` if conversion warnings rather than errors are acceptable. To avoid problems in the first place, do not use MySQL to check column content. It is safest (and often faster) to let the application ensure that it passes only valid values to the database.

With either of the strict mode options, you can cause errors to be treated as warnings by using `INSERT IGNORE` or `UPDATE IGNORE` rather than `INSERT` or `UPDATE` without `IGNORE`.

1.8.3.4 ENUM and SET Constraints

`ENUM` and `SET` columns provide an efficient way to define columns that can contain only a given set of values. See Section 11.4.4, “The ENUM Type”, and Section 11.4.5, “The SET Type”.

With strict mode enabled (see Section 5.1.8, “Server SQL Modes”), the definition of a `ENUM` or `SET` column acts as a constraint on values entered into the column. An error occurs for values that do not satisfy these conditions:

- An `ENUM` value must be one of those listed in the column definition, or the internal numeric equivalent thereof. The value cannot be the error value (that is, 0 or the empty string). For a column defined as `ENUM('a','b','c')`, values such as `''`, `'d'`, or `'ax'` are invalid and are rejected.
- A `SET` value must be the empty string or a value consisting only of the values listed in the column definition separated by commas. For a column defined as `SET('a','b','c')`, values such as `'d'` or `'a,b,c,d'` are invalid and are rejected.

Errors for invalid values can be suppressed in strict mode if you use `INSERT IGNORE` or `UPDATE IGNORE`. In this case, a warning is generated rather than an error. For `ENUM`, the value is inserted as the error member (`0`). For `SET`, the value is inserted as given except that any invalid substrings are deleted. For example, `'a,x,b,y'` results in a value of `'a,b'`.

1.9 Credits

The following sections list developers, contributors, and supporters that have helped to make MySQL what it is today.

1.9.1 Contributors to MySQL

Although Oracle Corporation and/or its affiliates own all copyrights in the `MySQL server` and the `MySQL manual`, we wish to recognize those who have made contributions of one kind or another to the `MySQL distribution`. Contributors are listed here, in somewhat random order:

- Gianmassimo Vigazzola `<qwerg@mbox.vol.it>` or `<qwerg@tin.it>`

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  Ported (and extended) the benchmark suite to `DBI`/`DBD`. Have been of great help with `crash-me` and running benchmarks. Some new date functions. The `mysql_setpermission` script.

- Alexis Mikhailov `<root@medinf.chuvashia.su>`

  User-defined functions (UDFs); `CREATE FUNCTION` and `DROP FUNCTION`.

- Andreas F. Bobak `<bobak@relog.ch>`

  The `AGGREGATE` extension to user-defined functions.

- Ross Wakelin `<R.Wakelin@march.co.uk>`

  Help to set up InstallShield for MySQL-Win32.

- Jethro Wright III `<jetman@li.net>`

  The `libmysql.dll` library.

- James Pereria `<jpereira@iafrica.com>`

  Mysqlmanager, a Win32 GUI tool for administering MySQL Servers.

- Curt Sampson <cjs@portal.ca>

  Porting of MIT-pthreads to NetBSD/Alpha and NetBSD 1.3/i386.

- Martin Ramsch <m.ramsch@computer.org>

  Examples in the MySQL Tutorial.

- Steve Harvey

  For making `mysqlaccess` more secure.

- Konark IA-64 Centre of Persistent Systems Private Limited

  Help with the Win64 port of the MySQL server.

- Albert Chin-A-Young.

  Configure updates for Tru64, large file support and better TCP wrappers support.

- John Birrell

  Emulation of `pthread_mutex()` for OS/2.

- Benjamin Pflugmann

  Extended `MERGE` tables to handle `INSERTS`. Active member on the MySQL mailing lists.

- Jocelyn Fournier

  Excellent spotting and reporting innumerable bugs (especially in the MySQL 4.1 subquery code).

- Marc Liyanage

  Maintaining the OS X packages and providing invaluable feedback on how to create OS X packages.

- Robert Rutherford

  Providing invaluable information and feedback about the QNX port.

- Previous developers of NDB Cluster

  Lots of people were involved in various ways summer students, master thesis students, employees. In total more than 100 people so too many to mention here. Notable name is Ataullah Dabaghi who up until 1999 contributed around a third of the code base. A special thanks also to developers of the AXE system which provided much of the architectural foundations for NDB Cluster with blocks, signals and crash tracing functionality. Also credit should be given to those who believed in the ideas enough to allocate of their budgets for its development from 1992 to present time.

- Google Inc.

  We wish to recognize Google Inc. for contributions to the MySQL distribution: Mark Callaghan's SMP Performance patches and other patches.

Other contributors, bugfinders, and testers: James H. Thompson, Maurizio Menghini, Wojciech Tryc, Luca Berra, Zarko Mocnik, Wim Bonis, Elmar Haneke, <jehamby@lightside>, <psmith@BayNetworks.com>, <duane@connect.com.au>, Ted Deppner <ted@psyber.com>, Mike Simons, Jaakko Hyvatti.

And lots of bug report/patches from the folks on the mailing list.

A big tribute goes to those that help us answer questions on the MySQL mailing lists:

- Daniel Koch <dkoch@amcity.com>

Irix setup.

- Luuk de Boer <luuk@wxs.nl>

  Benchmark questions.

- Tim Sailer <tps@users.buoy.com>

  DBD::mysql questions.

- Boyd Lynn Gerber <gerberb@zenez.com>

  SCO-related questions.

- Richard Mehalick <RM186061@shellus.com>

  xmysql-related questions and basic installation questions.

- Zeev Suraski <bourbon@netvision.net.il>

  Apache module configuration questions (log & auth), PHP-related questions, SQL syntax-related questions and other general questions.

- Francesc Guasch <frankie@citel.upc.es>

  General questions.

- Jonathan J Smith <jsmith@wtp.net>

  Questions pertaining to OS-specifics with Linux, SQL syntax, and other things that might need some work.

- David Sklar <sklar@student.net>

  Using MySQL from PHP and Perl.

- Alistair MacDonald <A.MacDonald@uel.ac.uk>

  Is flexible and can handle Linux and perhaps HP-UX.

- John Lyon <jlyon@imag.net>

  Questions about installing MySQL on Linux systems, using either .rpm files or compiling from source.

- Lorvid Ltd. <lorvid@WOLFENET.com>

  Simple billing/license/support/copyright issues.

- Patrick Sherrill <patrick@coconet.com>

  ODBC and VisualC++ interface questions.

- Randy Harmon <rjharmon@uptimecomputers.com>

  DBD, Linux, some SQL syntax questions.

1.9.2 Documenters and translators

The following people have helped us with writing the MySQL documentation and translating the documentation or error messages in MySQL.

- Paul DuBois

Ongoing help with making this manual correct and understandable. That includes rewriting Monty's and David's attempts at English into English as other people know it.

- Kim Aldale

  Helped to rewrite Monty's and David's early attempts at English into English.

- Michael J. Miller Jr. <mke@terrapin.turbolift.com>

  For the first MySQL manual. And a lot of spelling/language fixes for the FAQ (that turned into the MySQL manual a long time ago).

- Yan Cailin

  First translator of the MySQL Reference Manual into simplified Chinese in early 2000 on which the Big5 and HK coded versions were based.

- Jay Flaherty <fty@mediapulse.com>

  Big parts of the Perl DBI/DBD section in the manual.

- Paul Southworth <pauls@etext.org>, Ray Loyzaga <yar@cs.su.oz.au>

  Proof-reading of the Reference Manual.

- Therrien Gilbert <gilbert@ican.net>, Jean-Marc Pouyot <jmp@scalaire.fr>

  French error messages.

- Petr Snajdr, <snajdr@pvt.net>

  Czech error messages.

- Jaroslaw Lewandowski <jotel@itnet.com.pl>

  Polish error messages.

- Miguel Angel Fernandez Roiz

  Spanish error messages.

- Roy-Magne Mo <rmo@www.hivolda.no>

  Norwegian error messages and testing of MySQL 3.21.xx.

- Timur I. Bakeyev <root@timur.tatarstan.ru>

  Russian error messages.

- <brenno@dewinter.com> & Filippo Grassilli <phil@hyppo.com>

  Italian error messages.

- Dirk Munzinger <dirk@trinity.saar.de>

  German error messages.

- Billik Stefan <billik@sun.uniag.sk>

  Slovak error messages.

- Stefan Saroiu <tzoompy@cs.washington.edu>

  Romanian error messages.

- Peter Feher

  Hungarian error messages.

- Roberto M. Serqueira

  Portuguese error messages.

- Carsten H. Pedersen

  Danish error messages.

- Arjen Lentz

  Dutch error messages, completing earlier partial translation (also work on consistency and spelling).

1.9.3 Packages that support MySQL

The following is a list of creators/maintainers of some of the most important API/packages/applications that a lot of people use with MySQL.

We cannot list every possible package here because the list would then be way to hard to maintain. For other packages, please refer to the software portal at http://solutions.mysql.com/software/.

- Tim Bunce, Alligator Descartes

  For the `DBD` (Perl) interface.

- Andreas Koenig `<a.koenig@mind.de>`

  For the Perl interface for MySQL Server.

- Jochen Wiedmann `<wiedmann@neckar-alb.de>`

  For maintaining the Perl `DBD::mysql` module.

- Eugene Chan `<eugene@acenet.com.sg>`

  For porting PHP for MySQL Server.

- Georg Richter

  MySQL 4.1 testing and bug hunting. New PHP 5.0 `mysqli` extension (API) for use with MySQL 4.1 and up.

- Giovanni Maruzzelli `<maruzz@matrice.it>`

  For porting iODBC (Unix ODBC).

- Xavier Leroy `<Xavier.Leroy@inria.fr>`

  The author of LinuxThreads (used by the MySQL Server on Linux).

1.9.4 Tools that were used to create MySQL

The following is a list of some of the tools we have used to create MySQL. We use this to express our thanks to those that has created them as without these we could not have made MySQL what it is today.

- Free Software Foundation

  From whom we got an excellent compiler (`gcc`), an excellent debugger (`gdb` and the `libc` library (from which we have borrowed `strto.c` to get some code working in Linux).

- Free Software Foundation & The XEmacs development team

  For a really great editor/environment.

- Julian Seward

  Author of `valgrind`, an excellent memory checker tool that has helped us find a lot of otherwise hard to find bugs in MySQL.

- Dorothea Lütkehaus and Andreas Zeller

  For `DDD` (The Data Display Debugger) which is an excellent graphical front end to `gdb`).

1.9.5 Supporters of MySQL

Although Oracle Corporation and/or its affiliates own all copyrights in the `MySQL server` and the `MySQL manual`, we wish to recognize the following companies, which helped us finance the development of the `MySQL server`, such as by paying us for developing a new feature or giving us hardware for development of the `MySQL server`.

- VA Linux / Andover.net

  Funded replication.

- NuSphere

  Editing of the MySQL manual.

- Stork Design studio

  The MySQL Web site in use between 1998-2000.

- Intel

  Contributed to development on Windows and Linux platforms.

- Compaq

  Contributed to Development on Linux/Alpha.

- SWSoft

  Development on the embedded `mysqld` version.

- FutureQuest

  The `--skip-show-database` option.

Chapter 2 Installing and Upgrading MySQL

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This chapter describes how to obtain and install MySQL. A summary of the procedure follows and later sections provide the details. If you plan to upgrade an existing version of MySQL to a newer version rather than install MySQL for the first time, see Section 2.10.1, “Upgrading MySQL”, for information about upgrade procedures and about issues that you should consider before upgrading.

If you are interested in migrating to MySQL from another database system, see Section A.8, “MySQL 8.0 FAQ: Migration”, which contains answers to some common questions concerning migration issues.

Installation of MySQL generally follows the steps outlined here:

1. **Determine whether MySQL runs and is supported on your platform.**

   Please note that not all platforms are equally suitable for running MySQL, and that not all platforms on which MySQL is known to run are officially supported by Oracle Corporation. For information about those platforms that are officially supported, see http://www.mysql.com/support/supportedplatforms/database.html on the MySQL Web site.

2. **Choose which distribution to install.**

   Several versions of MySQL are available, and most are available in several distribution formats. You can choose from pre-packaged distributions containing binary (precompiled) programs or source code. When in doubt, use a binary distribution. Oracle also provides access to the MySQL source code for those who want to see recent developments and test new code. To determine which version and type of distribution you should use, see Section 2.1.1, “Which MySQL Version and Distribution to Install”.

3. **Download the distribution that you want to install.**

   For instructions, see Section 2.1.2, “How to Get MySQL”. To verify the integrity of the distribution, use the instructions in Section 2.1.3, “Verifying Package Integrity Using MD5 Checksums or GnuPG”.

4. **Install the distribution.**

   To install MySQL from a binary distribution, use the instructions in Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

   To install MySQL from a source distribution or from the current development source tree, use the instructions in Section 2.8, “Installing MySQL from Source”.

5. **Perform any necessary postinstallation setup.**

   After installing MySQL, see Section 2.9, “Postinstallation Setup and Testing” for information about making sure the MySQL server is working properly. Also refer to the information provided in Section 2.9.4, “Securing the Initial MySQL Account”. This section describes how to secure the initial MySQL `root` user account, *which has no password* until you assign one. The section applies whether you install MySQL using a binary or source distribution.

6. If you want to run the MySQL benchmark scripts, Perl support for MySQL must be available. See Section 2.11, “Perl Installation Notes”.

Instructions for installing MySQL on different platforms and environments is available on a platform by platform basis:

- **Unix, Linux, FreeBSD**

  For instructions on installing MySQL on most Linux and Unix platforms using a generic binary (for example, a `.tar.gz` package), see Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

  For information on building MySQL entirely from the source code distributions or the source code repositories, see Section 2.8, “Installing MySQL from Source”

  For specific platform help on installation, configuration, and building from source see the corresponding platform section:

  - Linux, including notes on distribution specific methods, see Section 2.5, “Installing MySQL on Linux”.
  - Solaris, including PKG and IPS formats, see Section 2.6, “Installing MySQL on Solaris”.
  - IBM AIX, see Section 2.6, “Installing MySQL on Solaris”.
  - FreeBSD, see Section 2.7, “Installing MySQL on FreeBSD”.

- **Microsoft Windows**

  For instructions on installing MySQL on Microsoft Windows, using either the MySQL Installer or Zipped binary, see Section 2.3, “Installing MySQL on Microsoft Windows”.

  For information about managing MySQL instances, see Section 2.3.4, “MySQL Notifier”.

  For details and instructions on building MySQL from source code using Microsoft Visual Studio, see Section 2.8, “Installing MySQL from Source”.

- **OS X**

  For installation on OS X, including using both the binary package and native PKG formats, see Section 2.4, “Installing MySQL on OS X”.

  For information on making use of an OS X Launch Daemon to automatically start and stop MySQL, see Section 2.4.3, “Installing a MySQL Launch Daemon”.

  For information on the MySQL Preference Pane, see Section 2.4.4, “Installing and Using the MySQL Preference Pane”.

2.1 General Installation Guidance

The immediately following sections contain the information necessary to choose, download, and verify your distribution. The instructions in later sections of the chapter describe how to install the distribution that you choose. For binary distributions, see the instructions at Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries” or the corresponding section for your platform if available. To build MySQL from source, use the instructions in Section 2.8, “Installing MySQL from Source”.

2.1.1 Which MySQL Version and Distribution to Install

MySQL is available on many operating systems and platforms. For information about platforms supported by GA releases of MySQL, see http://www.mysql.com/support/supportedplatforms/database.html. For development versions of MySQL, builds are available for a number of platforms at http://dev.mysql.com/downloads/mysql/8.0.html. To learn more about MySQL Support, see http://www.mysql.com/support/.

When preparing to install MySQL, decide which version and distribution format (binary or source) to use.

First, decide whether to install a development release or a General Availability (GA) release. Development releases have the newest features, but are not recommended for production use. GA releases, also called production or stable releases, are meant for production use. We recommend using the most recent GA release.

The naming scheme in MySQL 8.0 uses release names that consist of three numbers and an optional suffix; for example, **mysql-8.0.1-dmr**. The numbers within the release name are interpreted as follows:

- The first number (**8**) is the major version number.
- The second number (**0**) is the minor version number. Taken together, the major and minor numbers constitute the release series number. The series number describes the stable feature set.
- The third number (**1**) is the version number within the release series. This is incremented for each new bugfix release. In most cases, the most recent version within a series is the best choice.

Release names can also include a suffix to indicate the stability level of the release. Releases within a series progress through a set of suffixes to indicate how the stability level improves. The possible suffixes are:

- **dmr** indicates a development milestone release (DMR). MySQL development uses a milestone model, in which each milestone introduces a small subset of thoroughly tested features. From one milestone to the next, feature interfaces may change or features may even be removed, based on feedback provided by community members who try these earily releases. Features within milestone releases may be considered to be of pre-production quality.
- **rc** indicates a Release Candidate (RC). Release candidates are believed to be stable, having passed all of MySQL's internal testing. New features may still be introduced in RC releases, but the focus shifts to fixing bugs to stabilize features introduced earlier within the series.
- Absence of a suffix indicates a General Availability (GA) or Production release. GA releases are stable, having successfully passed through the earlier release stages, and are believed to be reliable, free of serious bugs, and suitable for use in production systems.

Development within a series begins with DMR releases, followed by RC releases, and finally reaches GA status releases.

After choosing which MySQL version to install, decide which distribution format to install for your operating system. For most use cases, a binary distribution is the right choice. Binary distributions are available in native format for many platforms, such as RPM packages for Linux or DMG packages for OS X. Distributions are also available in more generic formats such as Zip archives or compressed `tar` files. On Windows, you can use the MySQL Installer to install a binary distribution.

Under some circumstances, it may be preferable to install MySQL from a source distribution:

- You want to install MySQL at some explicit location. The standard binary distributions are ready to run at any installation location, but you might require even more flexibility to place MySQL components where you want.
- You want to configure `mysqld` with features that might not be included in the standard binary distributions. Here is a list of the most common extra options used to ensure feature availability:
  - `-DWITH_LIBWRAP=1` for TCP wrappers support.
  - `-DWITH_ZLIB={system|bundled}` for features that depend on compression
  - `-DWITH_DEBUG=1` for debugging support

  For additional information, see Section 2.8.4, “MySQL Source-Configuration Options”.
- You want to configure `mysqld` without some features that are included in the standard binary distributions.

- You want to read or modify the C and C++ code that makes up MySQL. For this purpose, obtain a source distribution.
- Source distributions contain more tests and examples than binary distributions.

2.1.2 How to Get MySQL

Check our downloads page at http://dev.mysql.com/downloads/ for information about the current version of MySQL and for downloading instructions. For a complete up-to-date list of MySQL download mirror sites, see http://dev.mysql.com/downloads/mirrors.html. You can also find information there about becoming a MySQL mirror site and how to report a bad or out-of-date mirror.

For RPM-based Linux platforms that use Yum as their package management system, MySQL can be installed using the MySQL Yum Repository. See Section 2.5.1, “Installing MySQL on Linux Using the MySQL Yum Repository” for details.

For Debian-based Linux platforms, MySQL can be installed using the MySQL APT Repository. See Section 2.5.2, “Installing MySQL on Linux Using the MySQL APT Repository” for details.

For SUSE Linux Enterprise Server (SLES) platforms, MySQL can be installed using the MySQL SLES Repository. See Section 2.5.3, “Installing MySQL on Linux Using the MySQL SLES Repository” for details.

To obtain the latest development source, see Section 2.8.3, “Installing MySQL Using a Development Source Tree”.

2.1.3 Verifying Package Integrity Using MD5 Checksums or GnuPG

After downloading the MySQL package that suits your needs and before attempting to install it, make sure that it is intact and has not been tampered with. There are three means of integrity checking:

- MD5 checksums
- Cryptographic signatures using `GnuPG`, the GNU Privacy Guard
- For RPM packages, the built-in RPM integrity verification mechanism

The following sections describe how to use these methods.

If you notice that the MD5 checksum or GPG signatures do not match, first try to download the respective package one more time, perhaps from another mirror site.

2.1.3.1 Verifying the MD5 Checksum

After you have downloaded a MySQL package, you should make sure that its MD5 checksum matches the one provided on the MySQL download pages. Each package has an individual checksum that you can verify against the package that you downloaded. The correct MD5 checksum is listed on the downloads page for each MySQL product, and you will compare it against the MD5 checksum of the file (product) that you download.

Each operating system and setup offers its own version of tools for checking the MD5 checksum. Typically the command is named `md5sum`, or it may be named `md5`, and some operating systems do not ship it at all. On Linux, it is part of the **GNU Text Utilities** package, which is available for a wide range of platforms. You can also download the source code from http://www.gnu.org/software/textutils/. If you have OpenSSL installed, you can use the command `openssl md5 package_name` instead. A Windows implementation of the `md5` command line utility is available from http://www.fourmilab.ch/md5/. `winMd5Sum` is a graphical MD5 checking tool that can be obtained from http://www.nullriver.com/index/products/winmd5sum. Our Microsoft Windows examples will assume the name `md5.exe`.

Linux and Microsoft Windows examples:

```
shell> md5sum mysql-standard-8.0.3-linux-i686.tar.gz
aaab65abbec64d5e907dcd41b8699945  mysql-standard-8.0.3-linux-i686.tar.gz
```

```
shell> md5.exe mysql-installer-community-8.0.3.msi
aaab65abbec64d5e907dcd41b8699945  mysql-installer-community-8.0.3.msi
```

You should verify that the resulting checksum (the string of hexadecimal digits) matches the one displayed on the download page immediately below the respective package.

Note

Make sure to verify the checksum of the *archive file* (for example, the `.zip`, `.tar.gz`, or `.msi` file) and not of the files that are contained inside of the archive. In other words, verify the file before extracting its contents.

2.1.3.2 Signature Checking Using GnuPG

Another method of verifying the integrity and authenticity of a package is to use cryptographic signatures. This is more reliable than using MD5 checksums, but requires more work.

We sign MySQL downloadable packages with GnuPG (GNU Privacy Guard). GnuPG is an Open Source alternative to the well-known Pretty Good Privacy (PGP) by Phil Zimmermann. Most Linux distributions ship with GnuPG installed by default. Otherwise, see http://www.gnupg.org/ for more information about GnuPG and how to obtain and install it.

To verify the signature for a specific package, you first need to obtain a copy of our public GPG build key, which you can download from http://pgp.mit.edu/. The key that you want to obtain is named `mysql-build@oss.oracle.com`. Alternatively, you can copy and paste the key directly from the following text:

```
-----BEGIN PGP PUBLIC KEY BLOCK-----
Version: GnuPG v1.4.5 (GNU/Linux)

mQGiBD4+owwRBAC14GIfUfCyEDSIePvEW3SAFUdJBtoQHH/nJKZyQT7h9bPlUWC3
RODjQReyCITRrdwyrKUGku2FmeVGwn2u2WmDMNABLnpprWPkBdCk96+OmSLN9brZ
fw2vOUgCmYv2hW0hyDHuvYlQA/BThQoADgj8AW6/0Lo7V1W9/8VuHP0gQwCgvzV3
BqOxRznNCRCRxAuAuVztHRcEAJooQK1+iSiunZMYD1WufeXfshc57S/+yeJkegNW
hxwR9pRWVArNYJdDRT+rf2RUe3vpquKNQU/hnEIUHJRQqYHo8gTxvxXNQc7fJYLV
K2HtkrPbP72vwsEKMYhhr0eKCbtLGfls9krjJ6sBgACyP/Vb7hiPwxh6rDZ7ITnE
kYpXBACmWpP8NJTkamEnPCia2ZoOHODANwpUkP43I7jsDmgtobZX9qnrAXw+uNDI
QJEXM6FSbi0LLtZciNlYsafwAPEOMDKpMqAK6IyisNtPvaLd8lH0bPAnWqcyefep
rv0sxxqUEMcM3o7wwgfN83POkDasDbs3pjwPhxvhz6//62zQJ7Q2TXlTUUwgUmVs
ZWFzZSBFbmdpbmVlcmluZyA8bXlzcWwtYnVpbGRAb3NzLm9yYWNsZS5jb20+iGwE
ExECACwCGyMCHgECF4ACGQEGCwkIBwMCBhUKCQgCAwUWAgMBAAUCWKcFIAUJHirJ
FAAKCRCMcY07UHLh9VcFAJ46pUyVd8BZ2r5CppMC1tmyQ3ceRgCfVPwuVsiS0VER
5WUqtAQDt+DoetCIaQQTEQIAKQIbIwYLCQgHAwIEFQIIAwQWAgMBAh4BAheAAhkB
BQJTAdRmBQkaZsvLAAoJEIxxjTtQcuH1X4MAoKNLWAbCBUj96637kv6Xa/fJuX5m
AJwPtmgDfjUe2iuhXdTrFEPT19SB6ohmBBMRAgAmAhsjBgsJCAcDAgQVAggDBBYC
AwECHgECF4AFAk53PioFCRP7AhUACgkQjHGNO1By4fUmzACeJdfqgc9gWTUhgmcM
AOmG4RjwuxcAoKfM+U8yMOGELi+TRif7MtKEms6piGkEExECACkCGyMGCwkIBwMC
BBUCCAMEFgIDAQIeAQIXgAIZAQUCUZSROgUJFTchqgAKCRCMcY07UHLh9YtAAJ9X
rA/ymlmozPZn+A9ls8/uwMcTsQCfaQMNq1dNkhH2kyByc3Rx9/W2xfqJARwEEAEC
AAYFAlAS6+UACgkQ8aIC+GoXHivrWwf/dtLk/x+NC2VMDlg+vOeM0qgG1IlhXZfi
NsEisvvGaz4m8fSFRGe+1bvvfDoKRhxiGXU48RusjixzvBb6KTMuY6JpOVfz9Dj3
H9spYriHa+i6rYySXZIpOhfLiMnTy7NH2OvYCyNzSS/ciIUACIfH/2NH8zNT5CNF
1uPNRs7HsHzzz7pOlTjtTWiF4cq/Ij6Z6CNrmdj+SiMvjYN9u6sdEKGtoNtpycgD
5HGKR+I7Nd/7v56yhaUe4FpuvsNXig86K9tI6MUFS8CUyy7Hj3kVBZOUWVBM053k
nGdALSygQr50DA3jMGKVl4ZnHje2RVWRmFTr5YWoRTMxUSQPMLpBNIkBHAQQAQIA
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DajKPpZv1RiZm8pMKqEb8WSw352xWoOcxuffjlsOEwvJ85SEGCAZ9tmIlkZOc7Ai
QONDvii9b8AYhQ60RIQC0HP2ASSmK0V92VeFPxHmAygdDQgZNVtbVxgnnt7oTNEu
VRXNY+z4OfBArp7R+cTsvijDRZY4kML1n22hUybwoxUEvjqZV2+JASIEEAECAAwF
AkrvOlQFAwASdQAACgkQlxC4m8pXrXxrPAgArXiNgZirNuBhfNCXlkzkCHLx5wnV
e4SmTpbWzTwWw7+qk7d4l9hlWtdImISORINzo7f4ShSUzJX2GciNaXhaHRo7+y5O
Zbu82jQb09aQQj/nibKYuqxqUrobTEm+DuYz3JUQZm2PsPcHLS8mX9cxvrJUncPG
nXEV0DRaq71SGWDprtkvBbp6i38aY3sIhYgz8wM5m1szKDtjywmBYcFehIdozt9z
hm7wZshzRWQX1+Rf/pIsnk+OzBIa34crSemTnacbV/B7278z2XAyziPNFuqz0xu+
iltOmYmayfNWAmumuw9NcuwWMlth6Mc2HLrpo0ZBheJ6iuDMPsHnwqdB/4kBIgQQ
AQIADAUCSwBd2gUDABJ1AAAKCRCXELibyletfP6tB/4m1w0BtlkJgtS6E+B/ns14
z4A4PGors+n+MYm05qzvi+EnDF/sytCmVcKeimrtvDcfoDtKAFFvJjcYXfnJdGWm
Pu0SJMRL5KKCirAKwZmU/saxOgoB5QLNw+DHPteJ3w9GmWlGxIqG1r15WC5duzBC
y3FsnjJYG3jaLnHOO9yXXb5h0kUTORfUKdvAr1gxF2KoatZWqGoaPPnHoqb88rjt
zk8I7gDqoXnzh8wLxa0ZYvfTC/McxdWTrwXLft+krmMQ18iIZEne2hvVLNJVuluU
oiWLeHA8iNCQ4W4WTdLc1mCnCjGTMX/MN41uLH0C9Ka4R6wEaqj4lPDk1B/1TV+Q
iQEiBBABAgAMBQJLEYGrBQMAEnUAAAoJEJcQuJvKV618naIH/2t9aH5mBTKBN6fU
qhrf79vIsjtI/QNS5qisBISZMX3/1/0Gu6WnxkPSfdCUJMWCjMcnVj7KU2wxTHHG
VpAStd9r2afUNxRyqZwzwyytktuZok0XngAEDYDDBS3ssu2R4uWLCsC2ysXEqO/5
tI5YrTWJZrfeIphTaYP5hxrMujvqy3kEwKKbiMz91cDeiLS+YCBcalj5n/1dMYf7
8U8C6ieurxAg/L8h6x25VM4Ilx4MmG2T8QGtkkUXd+Fd/KYWmf0LE5LLPknf0Hhw
oVslPXeinp4FsHK/5wzviv4YZpzuTqs9NlKcMsa4IuuPOB0FDf0pn+OFQbEg9QwY
2gCozK+JASIEEAECAAwFAksjTdQFAwASdQAACgkQlxC4m8pXrXwlogf/XBGbXRVX
LMaRN4SczOjwT3/tUCriTkb3v+zKjRG90zFhYAccjn7w+7jKQicjq6quQG1EH2X4
/Su6ps1lDLqGHHhiJW3ZhxQScLZmhdAYsh2qG4GP/UW3QjXG7c61t+H3olvWg2cr
wqCxxFZAgkAAkr9xcHWFZJEQeXoob6cCZObaUnHSANdmC6s5lUxXYa2bmL7Q3UB4
4KCzDvAfbPZKJOw9k0qb3lc11zx+vGdyZFbm4R0+3LPp/vT0b3GlSbbF9lU1GOXh
VaphrgFFa76dmjfHCkPplXAkK1VSIU/aPGAefduTFMdlSZpdMtJ5AULjGcszBDlR
pLlPxvqVa0ZpgIkBIgQQAQIADAUCSycmkgUDABJ1AAAKCRCXELibyletfHlNCACp
1YespiHfQt2alcscE5zgfETEHHic8Ai6pNkU9HT4TeWcFHEDe5QqfYcpjLrQvBXS
kSvxEittbyRdv+e+j5Z+HyHjiG8nAQBL6qy9eHqQE4+d7gYs6DTk7sG9ZMYphREb
ltzD+F4hVCQdLT8LNr0eVFN7ehqECScDaCG8/Qyti+l/0M902/Yn+mz0ilOiUdWJ
9x6LPaIINtb1gsYDEylLjwGIZmI0r5Kh9wYoV4vnNezFbxO1uRiW0B7iaPjIEsbt
OOKp7wx2aX+DM3N9F3BtaIY8XnzcnomNm83SNsgmgrZljpQltUnNqIhNM8DupQ+I
WOV5gtl6pTC7CgeVTVyRiQEiBBABAgAMBQJLOGXuBQMAEnUAAAoJEJcQuJvKV618
ll4IAKJ9mm4jb0c8fe9+uDI8eCJRbzNbVXm8zWzpA8GUtQAakwxoKv332QP1Wa1P
odni/e3EMhsSREOZJJv79YqGxGRBTE9Kb/VjM34nas4XSnXKW28XWhKyIw+XwQAi
nY2swFHh+83Htr/mwTdJfS2aEYl2zboBvd/JZCdhOGU2GH737S/3uEczoKkfVQ/w
OTM8X1xWwlYWqx23k/DsGcuDs9lA2g7Mx7DSqBtVjaTkn9h0zATzXLDkmP4SAUVj
cZ83WDpFre5WnizZjdXlBMM5OCexp5WpmzyHLTnaBFK4jEmnsk5C2Rnoyp8Ivz6g
Ecg1tRbEXijRw++d2TFYlJwLKtiJASIEEAECAAwFAktKMicFAwASdQAACgkQlxC4
m8pXrXxqHQgAuYY5scKrh0m/GS9EYnyC9494lOlO6iytU0CpE6oBC31M3hfX/Dbj
```

```
UbcS5szZNU+2CPYo4ujQLZ7suN7+tTjG6pZFfMevajT9+jsL+NPMF8RLdLOVYmbl
TmSQGNO+XGEYaKYH5oZIeIW5AKCgi2ozkdFlBBLAx7Kqo/FyybhkURFEcvEyVmgf
3KLV7IIiX/fYLfoCMCJ/Lcm9/llSFB1n8Nvg66Xd533DKoHjueD3jyaNAVlo2mq/
sIAv++kntvOiB3GDK5pfwHZ78WWiCpsWZpE5gzAnzJ1Y0WEigRo0PVLu3cLO0jLG
23d+H/CbfZ8rkajHJeCDQF7YVmP0t0nYpYkBIgQQAQIADAUCS1v+ZgUDABJ1AAAK
CRCXELibyletfNS/CACqt2TkB86mjqM+cJ74+dWBvJ2aFuURuxzm95i9Q/W/hU08
2iMbC3+0k2oD8CrTOe61P+3oRyLjv/UEDUNzLncNe2YsA9JeV+4hvPwH5Vp3Om13
089fCKZUbqslXNKkHiWYU+zAaZJXEuGRmRz0HbQIeAMOWF4oa226uo1e4ws1Jhc+
F3E/ApCRyFBqBUdL05hapQLditYpsBjIdiBGpjzidMLE2wX2W4ZpAdN0U6BIyIqR
mTPjbSkvzS9kSWFmfhQgnBDKEYJpVZgE1sN52rYC1sDeGeiuKxlzjVov9MMhYMWa
Zo3R5o3F2iIM/BK6FbC252lf/Mhu3ICuXujNBZNYiQEiBBABAgAMBQJLbSH4BQMA
EnUAAAoJEJcQuJvKV618kd0IAJLLwDH6gvgAlBFklQJXqQxUdcSOOVMAWtlHgWOy
ozjgomZZBkRL8dtCDr9YBMcj5czcQ3qpmLJdppXhKB+kJV2iUXfDMSFXwJ4wLfIs
8FNnXw8H5U01oBkGH/Ku6ngL9Vwt+MjYHtCWkw9QueUKZnDudX9qIzLAIt+mwSTu
A6+fY4VWIg40AA0v3exaQM55YR/UhlKunpGG9o8Qkq77dMEbTMpOmBoLbOMRB3Dd
MAvVU6G2l6Pcb7KobVCuOBnb6batXARV/G8sw+nzfJ16fr/KobZT2A6m+Jrqk4dl
F14ljLbz16O5JGUPAryN2G2ddBdSAy7dtFSVhWWiWC9n88q5Ag0EPj6jHRAIAO/h
iX8WzHWOMLJT54x/axeDdqn1rBDf5cWmaCWHN2ujNNlgpx5emoU9v7QStsNUCOGB
bXkeO4Ar7YG+jtSR33zqNh3y5kQ0YkY3dQ0wh6nsl+wh4XIIY/3TUZVtmdJeUBRH
JlfVNFYad2hX1guFI37Ny1PoZAFsxO82g+XB/Se8r/+sbmVcONdcdIeFKrE3FjLt
IjNQcxC6l9Q2Oy8KDxG/zvUZG3+H5i3tdRMyGgmuD6gEV0GXOHYUopzLeit1+Aa0
bCk36Mwbu+BeOw/CJW3+b0mB27hOaf9aCA855IP6fJFvtxcblq8nHIqhU3Dc9tec
sl9/S1xZ5S8ylG/xeRsAAwUH/i8KqmvAhq0X7DgCcYputwh37cuZlHOa1Ep07JRm
BCDgkdQXkGrsj2Wzw7Aw/TGdWWkmn2pxb8BRui5cfcZFO7c6vryi6FpJuLucX975
+eVY50ndWkPXkJ1HF4i+HJwRqE2zliN/RHMs4LJcwXQvvjD43EE3AO6eiVFbD+qA
AdxUFoOeLblKNBHPG7DPG9xL+Ni5rkE+TXShxsB7F0z7ZdJJZOG0JODmox7IstQT
GoaU9u41oyZTIiXPiFidJoIZCh7fdurP8pn3X+R5HUNXMr7M+ba8lSNxce/F3kmH
0L7rsKqdh9d/aVxhJINJ+inVDnrXWVoXu9GBjT8Nco1iU9SIVAQYEQIADAUCTnc9
7QUJE/sBuAASB2VHUEcAAQEJEIxxjTtQcuH1FJsAmwWK9vmwRJ/y9gTnJ8PWf0BV
roUTAKClYAhZuX2nUNwH4vlEJQHDqYa5yQ==
=HfUN
-----END PGP PUBLIC KEY BLOCK-----
```

To import the build key into your personal public GPG keyring, use `gpg --import`. For example, if you have saved the key in a file named `mysql_pubkey.asc`, the import command looks like this:

```
shell> gpg --import mysql_pubkey.asc
gpg: key 5072E1F5: public key "MySQL Release Engineering
<mysql-build@oss.oracle.com>" imported
gpg: Total number processed: 1
gpg:               imported: 1
gpg: no ultimately trusted keys found
```

You can also download the key from the public keyserver using the public key id, `5072E1F5`:

```
shell> gpg --recv-keys 5072E1F5
gpg: requesting key 5072E1F5 from hkp server keys.gnupg.net
gpg: key 5072E1F5: "MySQL Release Engineering <mysql-build@oss.oracle.com>"
1 new user ID
gpg: key 5072E1F5: "MySQL Release Engineering <mysql-build@oss.oracle.com>"
53 new signatures
gpg: no ultimately trusted keys found
gpg: Total number processed: 1
gpg:           new user IDs: 1
gpg:         new signatures: 53
```

If you want to import the key into your RPM configuration to validate RPM install packages, you should be able to import the key directly:

```
shell> rpm --import mysql_pubkey.asc
```

If you experience problems or require RPM specific information, see Section 2.1.3.4, “Signature Checking Using RPM”.

After you have downloaded and imported the public build key, download your desired MySQL package and the corresponding signature, which also is available from the download page. The signature file has the same name as the distribution file with an `.asc` extension, as shown by the examples in the following table.

Table 2.1 MySQL Package and Signature Files for Source files

File Type	File Name
Distribution file	`mysql-standard-8.0.3-linux-i686.tar.gz`
Signature file	`mysql-standard-8.0.3-linux-i686.tar.gz.asc`

Make sure that both files are stored in the same directory and then run the following command to verify the signature for the distribution file:

```
shell> gpg --verify package_name.asc
```

If the downloaded package is valid, you will see a "Good signature" similar to:

```
shell> gpg --verify mysql-standard-8.0.3-linux-i686.tar.gz.asc
gpg: Signature made Tue 01 Feb 2011 02:38:30 AM CST using DSA key ID 5072E1F5
gpg: Good signature from "MySQL Release Engineering <mysql-build@oss.oracle.com>"
```

The `Good signature` message indicates that the file signature is valid, when compared to the signature listed on our site. But you might also see warnings, like so:

```
shell> gpg --verify mysql-standard-8.0.3-linux-i686.tar.gz.asc
gpg: Signature made Wed 23 Jan 2013 02:25:45 AM PST using DSA key ID 5072E1F5
gpg: checking the trustdb
gpg: no ultimately trusted keys found
gpg: Good signature from "MySQL Release Engineering <mysql-build@oss.oracle.com>"
gpg: WARNING: This key is not certified with a trusted signature!
gpg:          There is no indication that the signature belongs to the owner.
Primary key fingerprint: A4A9 4068 76FC BD3C 4567  70C8 8C71 8D3B 5072 E1F5
```

That is normal, as they depend on your setup and configuration. Here are explanations for these warnings:

- *gpg: no ultimately trusted keys found*: This means that the specific key is not "ultimately trusted" by you or your web of trust, which is okay for the purposes of verifying file signatures.
- *WARNING: This key is not certified with a trusted signature! There is no indication that the signature belongs to the owner.*: This refers to your level of trust in your belief that you possess our real public key. This is a personal decision. Ideally, a MySQL developer would hand you the key in person, but more commonly, you downloaded it. Was the download tampered with? Probably not, but this decision is up to you. Setting up a web of trust is one method for trusting them.

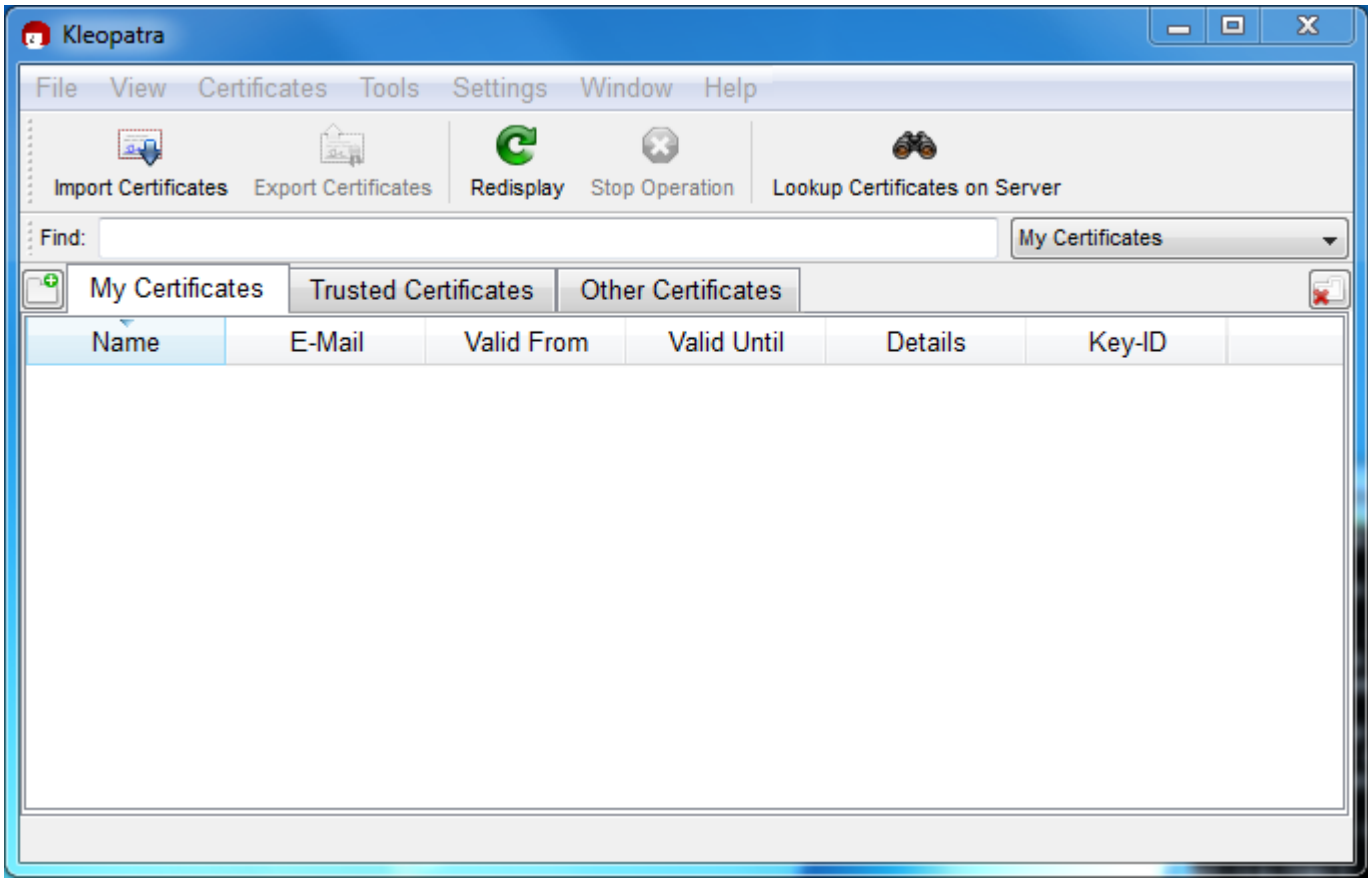
See the GPG documentation for more information on how to work with public keys.

2.1.3.3 Signature Checking Using Gpg4win for Windows

The Section 2.1.3.2, “Signature Checking Using GnuPG” section describes how to verify MySQL downloads using GPG. That guide also applies to Microsoft Windows, but another option is to use a GUI tool like Gpg4win. You may use a different tool but our examples are based on Gpg4win, and utilize its bundled `Kleopatra` GUI.

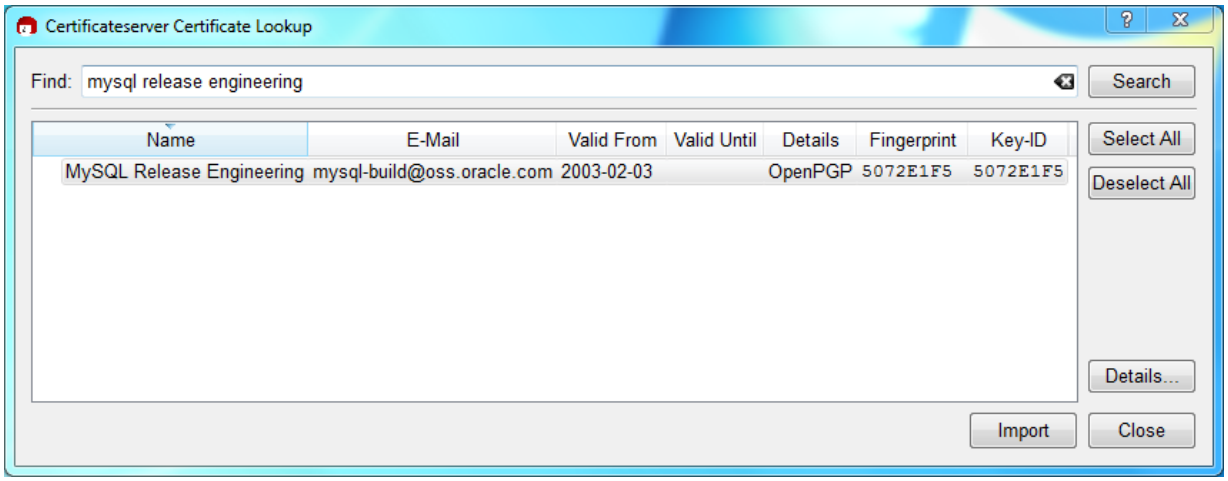
Download and install Gpg4win, and then load Kleopatra. The dialog should look similar to:

Figure 2.1 Initial screen after loading Kleopatra



Next, add the MySQL Release Engineering certificate. Do this by clicking **File**, **Lookup Certificates on Server**. Type "Mysql Release Engineering" into the search box and press **Search**.

Figure 2.2 Finding the MySQL Release Engineering certificate

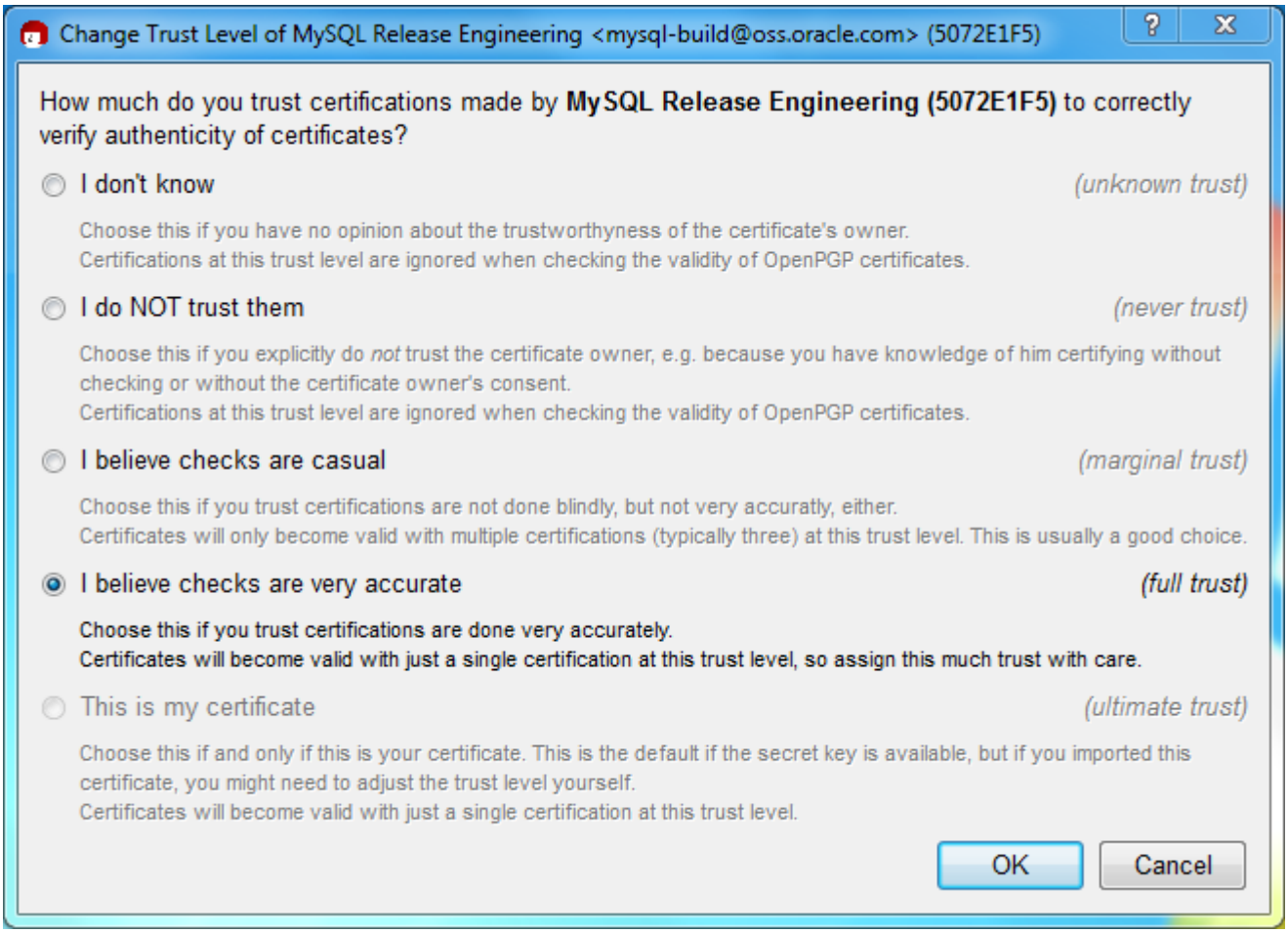


Select the "MySQL Release Engineering" certificate. The Fingerprint and Key-ID must be "5072E1F5", or choose **Details...** to confirm the certificate is valid. Now, import it by clicking **Import**. An import dialog will be displayed, choose **Okay**, and this certificate will now be listed under the **Imported Certificates** tab.

Next, configure the trust level for our certificate. Select our certificate, then from the main menu select **Certificates**, **Change Owner Trust...**. We suggest choosing **I believe checks are very accurate** for

our certificate, as otherwise you might not be able to verify our signature. Select **I believe checks are very accurate** and then press **OK**.

Figure 2.3 Changing the Trust level



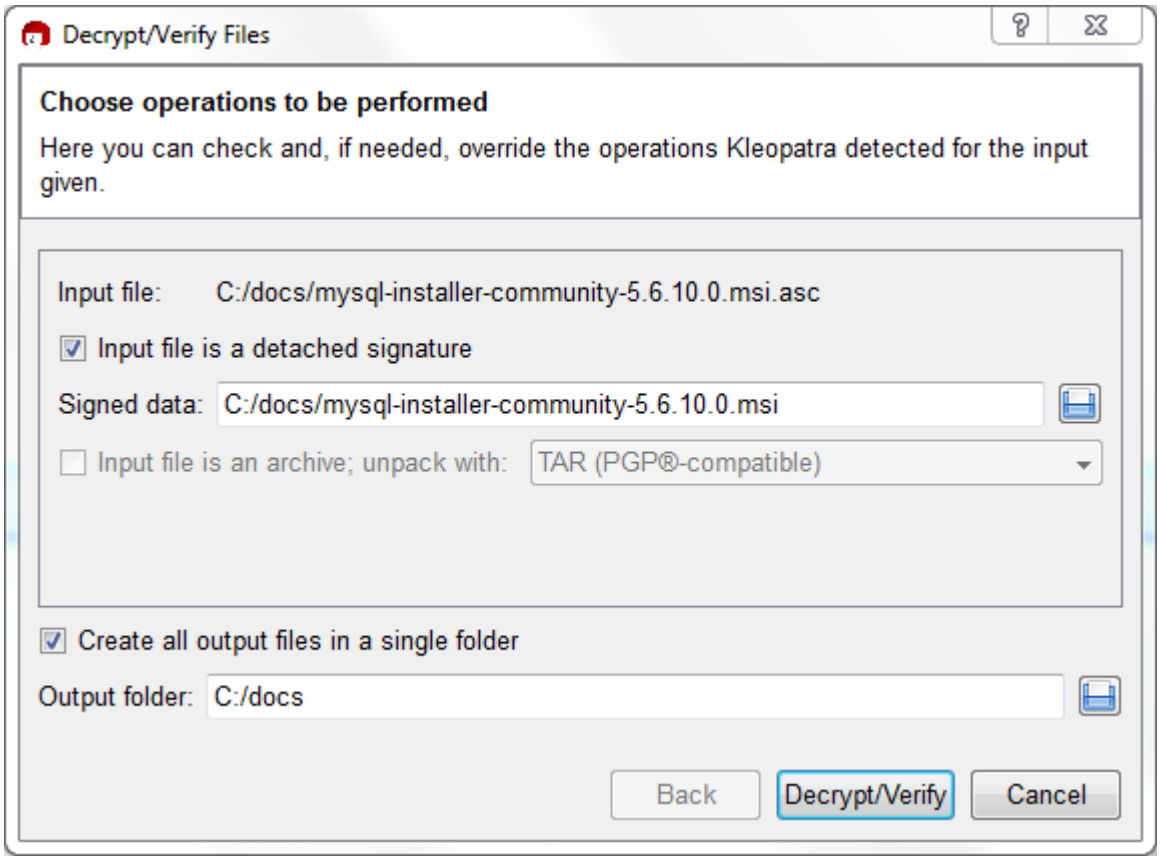
Next, verify the downloaded MySQL package file. This requires files for both the packaged file, and the signature. The signature file must have the same name as the packaged file but with an appended `.asc` extension, as shown by the example in the following table. The signature is linked to on the downloads page for each MySQL product. You must create the `.asc` file with this signature.

Table 2.2 MySQL Package and Signature Files for MySQL Installer for Microsoft Windows

File Type	File Name
Distribution file	`mysql-installer-community-8.0.3.msi`
Signature file	`mysql-installer-community-8.0.3.msi.asc`

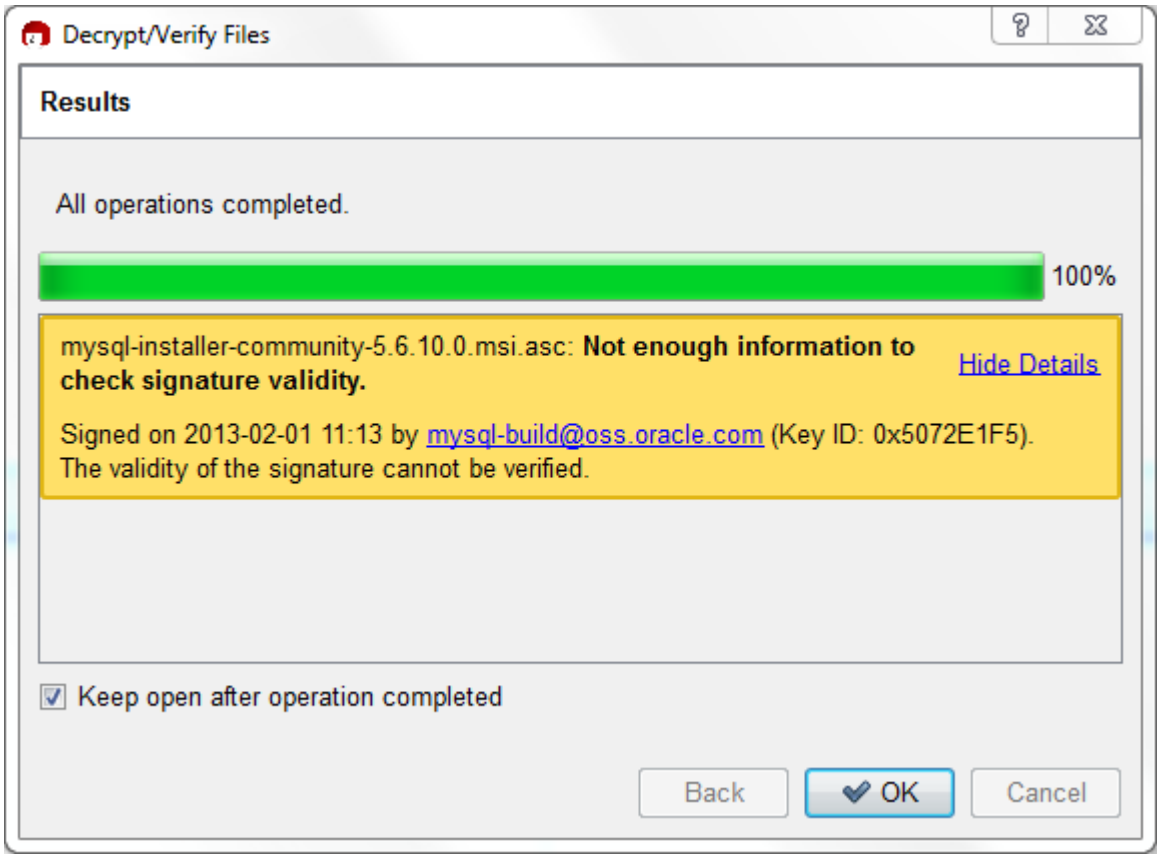
Make sure that both files are stored in the same directory and then run the following command to verify the signature for the distribution file. Either drag and drop the signature (`.asc`) file into Kleopatra, or load the dialog from **File**, **Decrypt/Verify Files...**, and then choose either the `.msi` or `.asc` file.

Figure 2.4 The Decrypt/Verify Files dialog



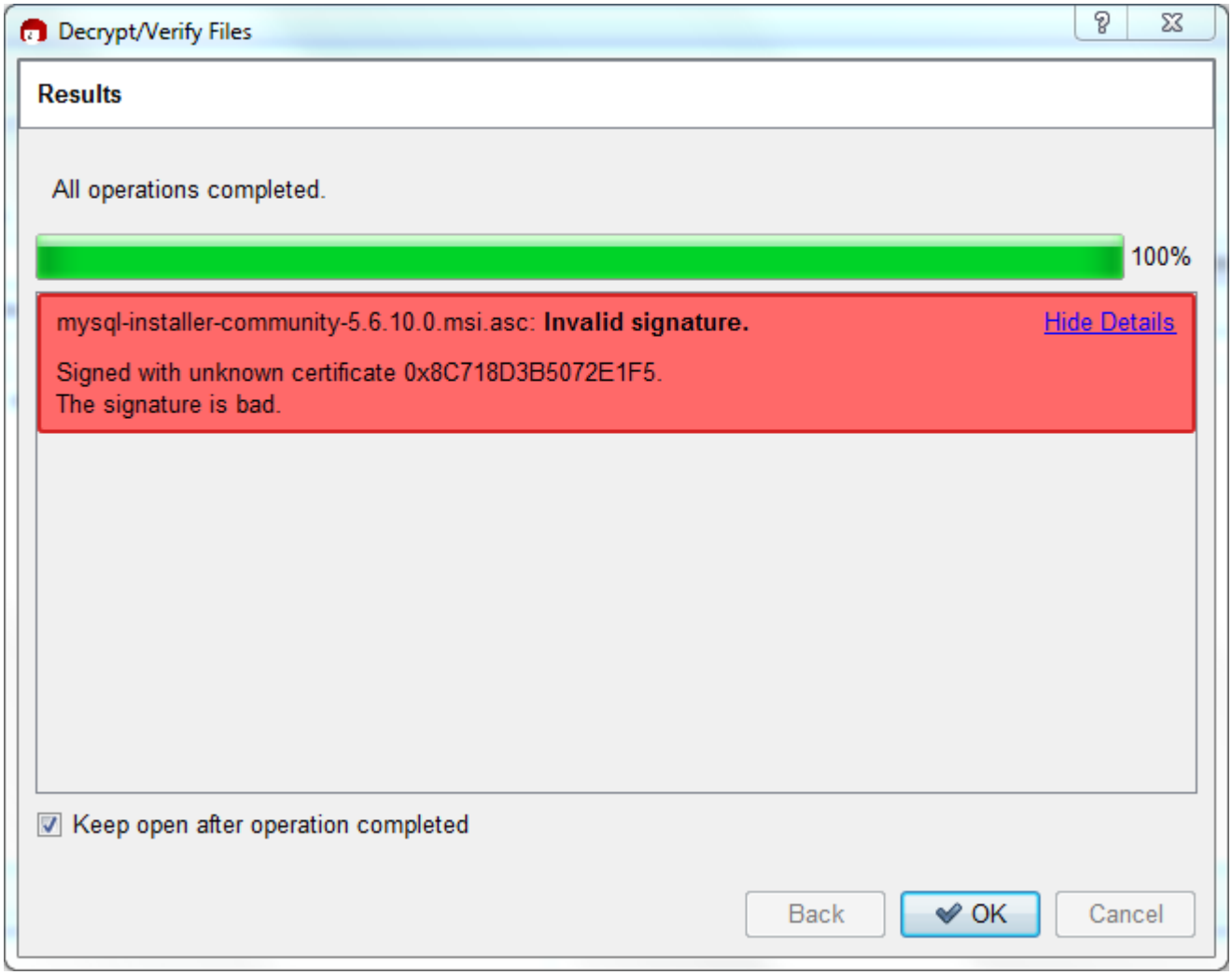
Click **Decrypt/Verify** to check the file. The two most common results will look like the following, and although the yellow warning looks problematic, the following means that the file check passed with success. You may now run this installer.

Figure 2.5 The Decrypt/Verify Results: Good



Seeing a red "The signature is bad" error means the file is invalid. Do not execute the MSI file if you see this error.

Figure 2.6 The Decrypt/Verify Results: Bad



The Section 2.1.3.2, “Signature Checking Using GnuPG” section explains why you probably don't see a green `Good signature` result.

2.1.3.4 Signature Checking Using RPM

For RPM packages, there is no separate signature. RPM packages have a built-in GPG signature and MD5 checksum. You can verify a package by running the following command:

```
shell> rpm --checksig package_name.rpm
```

Example:

```
shell> rpm --checksig MySQL-server-8.0.3-0.linux_glibc2.5.i386.rpm
MySQL-server-8.0.3-0.linux_glibc2.5.i386.rpm: md5 gpg OK
```

Note

If you are using RPM 4.1 and it complains about `(GPG) NOT OK (MISSING KEYS: GPG#5072e1f5)`, even though you have imported the MySQL public build key into your own GPG keyring, you need to import the key into the RPM keyring first. RPM 4.1 no longer uses your personal GPG keyring (or GPG itself). Rather, RPM maintains a separate keyring because it is a system-wide

application and a user's GPG public keyring is a user-specific file. To import the MySQL public key into the RPM keyring, first obtain the key, then use `rpm --import` to import the key. For example:

```
shell> gpg --export -a 5072e1f5 > 5072e1f5.asc
shell> rpm --import 5072e1f5.asc
```

Alternatively, `rpm` also supports loading the key directly from a URL, and you can use this manual page:

```
shell> rpm --import http://dev.mysql.com/doc/refman/8.0/en/checking-gpg-signature.html
```

If you need to obtain the MySQL public key, see Section 2.1.3.2, “Signature Checking Using GnuPG”.

2.1.4 Installation Layouts

The installation layout differs for different installation types (for example, native packages, binary tarballs, and source tarballs), which can lead to confusion when managing different systems or using different installation sources. The individual layouts are given in the corresponding installation type or platform chapter, as described following. Note that the layout of installations from vendors other than Oracle may differ from these layouts.

- Section 2.3.1, “MySQL Installation Layout on Microsoft Windows”
- Section 2.8.1, “MySQL Layout for Source Installation”
- Table 2.3, “MySQL Installation Layout for Generic Unix/Linux Binary Package”
- Table 2.10, “MySQL Installation Layout for Linux RPM Packages from the MySQL Developer Zone”
- Table 2.6, “MySQL Installation Layout on OS X”

2.1.5 Compiler-Specific Build Characteristics

In some cases, the compiler used to build MySQL affects the features available for use. The notes in this section apply for binary distributions provided by Oracle Corporation or that you compile yourself from source.

`icc` (Intel C++ Compiler) Builds

A server built with `icc` has these characteristics:

- SSL support is not included.

2.2 Installing MySQL on Unix/Linux Using Generic Binaries

Oracle provides a set of binary distributions of MySQL. These include generic binary distributions in the form of compressed `tar` files (files with a `.tar.gz` extension) for a number of platforms, and binaries in platform-specific package formats for selected platforms.

This section covers the installation of MySQL from a compressed `tar` file binary distribution. For other platform-specific package formats, see the other platform-specific sections. For example, for Windows distributions, see Section 2.3, “Installing MySQL on Microsoft Windows”.

To obtain MySQL, see Section 2.1.2, “How to Get MySQL”.

MySQL compressed `tar` file binary distributions have names of the form `mysql-`*`VERSION`*`-`*`OS`*`.tar.gz`, where *`VERSION`* is a number (for example, `8.0.3`), and *`OS`* indicates

the type of operating system for which the distribution is intended (for example, `pc-linux-i686` or `winx64`).

Warning

If you have previously installed MySQL using your operating system native package management system, such as `yum` or `apt-get`, you may experience problems installing using a native binary. Make sure your previous MySQL installation has been removed entirely (using your package management system), and that any additional files, such as old versions of your data files, have also been removed. You should also check for configuration files such as `/etc/my.cnf` or the `/etc/mysql` directory and delete them.

For information about replacing third-party packages with official MySQL packages, see the related Apt guide or Yum guide.

Warning

MySQL has a dependency on the `libaio` library. Data directory initialization and subsequent server startup steps will fail if this library is not installed locally. If necessary, install it using the appropriate package manager. For example, on Yum-based systems:

```
shell> yum search libaio  # search for info
shell> yum install libaio # install library
```

Or, on APT-based systems:

```
shell> apt-cache search libaio # search for info
shell> apt-get install libaio1 # install library
```

If you run into problems and need to file a bug report, please use the instructions in Section 1.7, “How to Report Bugs or Problems”.

On Unix, to install a compressed `tar` file binary distribution, unpack it at the installation location you choose (typically `/usr/local/mysql`). This creates the directories shown in the following table.

Table 2.3 MySQL Installation Layout for Generic Unix/Linux Binary Package

Directory	Contents of Directory
`bin`	`mysqld` server, client and utility programs
`data`	Log files, databases
`docs`	MySQL manual in Info format
`man`	Unix manual pages
`include`	Include (header) files
`lib`	Libraries
`share`	Miscellaneous support files, including error messages, sample configuration files, SQL for database installation

Debug versions of the `mysqld` binary are available as `mysqld-debug`. To compile your own debug version of MySQL from a source distribution, use the appropriate configuration options to enable debugging support. See Section 2.8, “Installing MySQL from Source”.

To install and use a MySQL binary distribution, the command sequence looks like this:

```
shell> groupadd mysql
shell> useradd -r -g mysql -s /bin/false mysql
shell> cd /usr/local
shell> tar zxvf /path/to/mysql-VERSION-OS.tar.gz
shell> ln -s full-path-to-mysql-VERSION-OS mysql
shell> cd mysql
shell> mkdir mysql-files
shell> chmod 750 mysql-files
shell> chown -R mysql .
shell> chgrp -R mysql .
shell> bin/mysqld --initialize --user=mysql
shell> bin/mysql_ssl_rsa_setup
shell> chown -R root .
shell> chown -R mysql data mysql-files
shell> bin/mysqld_safe --user=mysql &
# Next command is optional
shell> cp support-files/mysql.server /etc/init.d/mysql.server
```

Note

This procedure assumes that you have `root` (administrator) access to your system. Alternatively, you can prefix each command using the `sudo` (Linux) or `pfexec` (Solaris) command.

The `mysql-files` directory provides a convenient location to use as the value of the `secure_file_priv` system variable that limits import/export operations to a specific directory. See Section 5.1.5, “Server System Variables”.

A more detailed version of the preceding description for installing a binary distribution follows.

Create a mysql User and Group

If your system does not already have a user and group to use for running `mysqld`, you may need to create one. The following commands add the `mysql` group and the `mysql` user. You might want to call the user and group something else instead of `mysql`. If so, substitute the appropriate name in the following instructions. The syntax for `useradd` and `groupadd` may differ slightly on different versions of Unix, or they may have different names such as `adduser` and `addgroup`.

```
shell> groupadd mysql
shell> useradd -r -g mysql -s /bin/false mysql
```

Note

Because the user is required only for ownership purposes, not login purposes, the `useradd` command uses the `-r` and `-s /bin/false` options to create a user that does not have login permissions to your server host. Omit these options if your `useradd` does not support them.

Obtain and Unpack the Distribution

Pick the directory under which you want to unpack the distribution and change location into it. The example here unpacks the distribution under `/usr/local`. The instructions, therefore, assume that you have permission to create files and directories in `/usr/local`. If that directory is protected, you must perform the installation as `root`.

```
shell> cd /usr/local
```

Obtain a distribution file using the instructions in Section 2.1.2, “How to Get MySQL”. For a given release, binary distributions for all platforms are built from the same MySQL source distribution.

Unpack the distribution, which creates the installation directory. `tar` can uncompress and unpack the distribution if it has `z` option support:

```
shell> tar zxvf /path/to/mysql-VERSION-OS.tar.gz
```

The `tar` command creates a directory named `mysql-VERSION-OS`.

To install MySQL from a compressed `tar` file binary distribution, your system must have GNU `gunzip` to uncompress the distribution and a reasonable `tar` to unpack it. If your `tar` program supports the `z` option, it can both uncompress and unpack the file.

GNU `tar` is known to work. The standard `tar` provided with some operating systems is not able to unpack the long file names in the MySQL distribution. You should download and install GNU `tar`, or if available, use a preinstalled version of GNU tar. Usually this is available as `gnutar`, `gtar`, or as `tar` within a GNU or Free Software directory, such as `/usr/sfw/bin` or `/usr/local/bin`. GNU `tar` is available from http://www.gnu.org/software/tar/.

If your `tar` does not have `z` option support, use `gunzip` to unpack the distribution and `tar` to unpack it. Replace the preceding `tar` command with the following alternative command to uncompress and extract the distribution:

```
shell> gunzip < /path/to/mysql-VERSION-OS.tar.gz | tar xvf -
```

Next, create a symbolic link to the installation directory created by `tar`:

```
shell> ln -s full-path-to-mysql-VERSION-OS mysql
```

The `ln` command makes a symbolic link to the installation directory. This enables you to refer more easily to it as `/usr/local/mysql`. To avoid having to type the path name of client programs always when you are working with MySQL, you can add the `/usr/local/mysql/bin` directory to your `PATH` variable:

```
shell> export PATH=$PATH:/usr/local/mysql/bin
```

Perform Postinstallation Setup

The remainder of the installation process involves setting distribution ownership and access permissions, initializing the data directory, starting the MySQL server, and setting up the configuration file. For instructions, see Section 2.9, “Postinstallation Setup and Testing”.

2.3 Installing MySQL on Microsoft Windows

Important

MySQL Community 8.0 Server requires the Microsoft Visual C++ 2015 Redistributable Package to run on Windows platforms. Users should make sure the package has been installed on the system before installing the server. The package is available at the Microsoft Download Center.

MySQL is available for Microsoft Windows, for both 32-bit and 64-bit versions. For supported Windows platform information, see http://www.mysql.com/support/supportedplatforms/database.html.

There are different methods to install MySQL on Microsoft Windows.

MySQL Installer Method

The simplest and recommended method is to download MySQL Installer (for Windows) and let it install and configure all of the MySQL products on your system. Here is how:

1. Download MySQL Installer from http://dev.mysql.com/downloads/installer/ and execute it.

Note

Unlike the standard MySQL Installer, the smaller "web-community" version does not bundle any MySQL applications but it will download the MySQL products you choose to install.

2. Choose the appropriate **Setup Type** for your system. Typically you will choose **Developer Default** to install MySQL server and other MySQL tools related to MySQL development, helpful tools like MySQL Workbench. Or, choose the **Custom** setup type to manually select your desired MySQL products.

Note

Multiple versions of MySQL server can exist on a single system. You can choose one or multiple versions.

3. Complete the installation process by following the MySQL Installation wizard's instructions. This will install several MySQL products and start the MySQL server.

MySQL is now installed. If you configured MySQL as a service, then Windows will automatically start MySQL server every time you restart your system.

Note

You probably also installed other helpful MySQL products like MySQL Workbench and MySQL Notifier on your system. Consider loading Chapter 30, *MySQL Workbench* to check your new MySQL server connection, and Section 2.3.4, “MySQL Notifier” to view the connection's status. By default, these two programs automatically start after installing MySQL.

This process also installs the MySQL Installer application on your system, and later you can use MySQL Installer to upgrade or reconfigure your MySQL products.

Additional Installation Information

It is possible to run MySQL as a standard application or as a Windows service. By using a service, you can monitor and control the operation of the server through the standard Windows service management tools. For more information, see Section 2.3.5.8, “Starting MySQL as a Windows Service”.

Generally, you should install MySQL on Windows using an account that has administrator rights. Otherwise, you may encounter problems with certain operations such as editing the `PATH` environment variable or accessing the `Service Control Manager`. Once installed, MySQL does not need to be executed using a user with Administrator privileges.

For a list of limitations on the use of MySQL on the Windows platform, see Section C.10.5, “Windows Platform Limitations”.

In addition to the MySQL Server package, you may need or want additional components to use MySQL with your application or development environment. These include, but are not limited to:

- To connect to the MySQL server using ODBC, you must have a Connector/ODBC driver. For more information, including installation and configuration instructions, see MySQL Connector/ODBC Developer Guide.

Note

MySQL Installer will install and configure Connector/ODBC for you.

- To use MySQL server with .NET applications, you must have the Connector/Net driver. For more information, including installation and configuration instructions, see MySQL Connector/Net Developer Guide.

  Note

  MySQL Installer will install and configure MySQL Connector/Net for you.

MySQL distributions for Windows can be downloaded from http://dev.mysql.com/downloads/. See Section 2.1.2, “How to Get MySQL”.

MySQL for Windows is available in several distribution formats, detailed here. Generally speaking, you should use MySQL Installer. It contains more features and MySQL products than the older MSI, is simpler to use than the Zip file, and you need no additional tools to get MySQL up and running. MySQL Installer automatically installs MySQL Server and additional MySQL products, creates an options file, starts the server, and enables you to create default user accounts. For more information on choosing a package, see Section 2.3.2, “Choosing An Installation Package”.

- A MySQL Installer distribution includes MySQL Server and additional MySQL products including MySQL Workbench, MySQL Notifier, and MySQL for Excel. MySQL Installer can also be used to upgrade these products in the future.

  For instructions on installing MySQL using MySQL Installer, see Section 2.3.3, “MySQL Installer for Windows”.

- The standard binary distribution (packaged as a Zip file) contains all of the necessary files that you unpack into your chosen location. This package contains all of the files in the full Windows MSI Installer package, but does not include an installation program.

  For instructions on installing MySQL using the Zip file, see Section 2.3.5, “Installing MySQL on Microsoft Windows Using a noinstall Zip Archive”.

- The source distribution format contains all the code and support files for building the executables using the Visual Studio compiler system.

  For instructions on building MySQL from source on Windows, see Section 2.8, “Installing MySQL from Source”.

MySQL on Windows considerations:

- **Large Table Support**

  If you need tables with a size larger than 4GB, install MySQL on an NTFS or newer file system. Do not forget to use `MAX_ROWS` and `AVG_ROW_LENGTH` when you create tables. See Section 13.1.15, “CREATE TABLE Syntax”.

- **MySQL and Virus Checking Software**

  Virus-scanning software such as Norton/Symantec Anti-Virus on directories containing MySQL data and temporary tables can cause issues, both in terms of the performance of MySQL and the virus-scanning software misidentifying the contents of the files as containing spam. This is due to the fingerprinting mechanism used by the virus-scanning software, and the way in which MySQL rapidly updates different files, which may be identified as a potential security risk.

  After installing MySQL Server, it is recommended that you disable virus scanning on the main directory (`datadir`) used to store your MySQL table data. There is usually a system built into the virus-scanning software to enable specific directories to be ignored.

  In addition, by default, MySQL creates temporary files in the standard Windows temporary directory. To prevent the temporary files also being scanned, configure a separate temporary directory for

MySQL temporary files and add this directory to the virus scanning exclusion list. To do this, add a configuration option for the `tmpdir` parameter to your `my.ini` configuration file. For more information, see Section 2.3.5.2, “Creating an Option File”.

2.3.1 MySQL Installation Layout on Microsoft Windows

For MySQL 8.0 on Windows, the default installation directory is `C:\Program Files\MySQL\MySQL Server 8.0`. Some Windows users prefer to install in `C:\mysql`, the directory that formerly was used as the default. However, the layout of the subdirectories remains the same.

All of the files are located within this parent directory, using the structure shown in the following table.

Table 2.4 Default MySQL Installation Layout for Microsoft Windows

Directory	Contents of Directory	Notes
`bin`	`mysqld` server, client and utility programs	
`%PROGRAMDATA%\MySQL\MySQL Server 8.0\`	Log files, databases	The Windows system variable `%PROGRAMDATA%` defaults to `C:\ProgramData`
`examples`	Example programs and scripts	
`include`	Include (header) files	
`lib`	Libraries	
`share`	Miscellaneous support files, including error messages, character set files, sample configuration files, SQL for database installation	

If you install MySQL using the MySQL Installer, this package creates and sets up the data directory that the installed server will use, and also creates a pristine “template” data directory named `data` under the installation directory. After an installation has been performed using this package, the template data directory can be copied to set up additional MySQL instances. See Section 5.7, “Running Multiple MySQL Instances on One Machine”.

2.3.2 Choosing An Installation Package

For MySQL 8.0, there are multiple installation package formats to choose from when installing MySQL on Windows.

Note

Program Database (PDB) files (with file name extension `pdb`) provide information for debugging your MySQL installation in the event of a problem. These files are included in ZIP Archive distributions (but not MSI distributions) of MySQL.

- **MySQL Installer**: This package has a file name similar to `mysql-installer-community-8.0.3.0.msi` or `mysql-installer-commercial-8.0.3.0.msi`, and utilizes MSIs to automatically install MySQL server and other products. It will download and apply updates to itself, and for each of the installed products. It also configures the additional non-server products.

  The installed products are configurable, and this includes: documentation with samples and examples, connectors (such as C, C++, J, NET, and ODBC), MySQL Workbench, MySQL Notifier, MySQL for Excel, and the MySQL Server with its components.

Note

MySQL Installer does not include debugging binaries/information components (including PDB files). These are available in a separate Zip archive named `mysql-VERSION-winx64-debug-test.zip` for 64-bit and `mysql-VERSION-win32-debug-test.zip` for 32-bit.

MySQL Installer operates on all MySQL supported versions of Windows (see http://www.mysql.com/support/supportedplatforms/database.html).

Note

Because MySQL Installer is not a native component of Microsoft Windows and depends on .NET, it will not work on minimal installation options like the "Server Core" version of Windows Server 2008.

For instructions on installing MySQL using MySQL Installer, see Section 2.3.3, “MySQL Installer for Windows”.

- **The Noinstall Archives**: These packages contain the files found in the complete installation package, with the exception of the GUI. This format does not include an automated installer, and must be manually installed and configured.

Note

The noinstall archives are split into two separate Zip files. The main package is named `mysql-VERSION-winx64.zip` for 64-bit and `mysql-VERSION-win32.zip` for 32-bit. This contains the components needed to use MySQL on your system. The optional MySQL test suite, MySQL benchmark suite, and debugging binaries/information components (including PDB files) are in a separate Zip file named `mysql-VERSION-winx64-debug-test.zip` for 64-bit and `mysql-VERSION-win32-debug-test.zip` for 32-bit.

MySQL Installer is recommended for most users.

Your choice of install package affects the installation process you must follow. If you choose to use MySQL Installer, see Section 2.3.3, “MySQL Installer for Windows”. If you choose to install a Noinstall archive, see Section 2.3.5, “Installing MySQL on Microsoft Windows Using a noinstall Zip Archive”.

2.3.3 MySQL Installer for Windows

MySQL Installer is a standalone application designed to ease the complexity of installing and managing MySQL products that run on Microsoft Windows. It supports the following MySQL products:

- MySQL Servers

  A single MySQL Installer instance can install and manage multiple MySQL server versions. For example, a single MySQL Installer instance can install (and update) versions 5.6, 5.7, and 8.0 on the same host. A host *cannot* have both community and commercial editions of MySQL server installed.

- MySQL Applications

  MySQL Workbench, MySQL Shell, MySQL Router, MySQL for Visual Studio, MySQL for Excel, MySQL Notifier, and MySQL Utilities.

- MySQL Connectors

  MySQL Connector/Net, MySQL Connector/Python, Connector/ODBC, MySQL Connector/J, Connector/C, and Connector/C++.

- Documentation and Samples

MySQL Reference Manuals (by version) in PDF format and MySQL database samples (by version).

MySQL Installer Community Edition

Download this edition from http://dev.mysql.com/downloads/installer/ to install the community version of all MySQL products for Windows. Select one of the following MySQL Installer package options:

- *Web:* Contains MySQL Installer and configuration files only. The web package downloads only the MySQL products you select to install, but it requires an internet connection for each download. The size of this file is approximately 2 MB; the name of the file has the form `mysql-installer-community-web-VERSION.N.msi` where `VERSION` is the MySQL server version number such as 5.7 and `N` is the package number, which begins at 0.
- *Full:* Bundles all of the MySQL products for Windows (including the MySQL server). The file size is over 300 MB, and its name has the form `mysql-installer-community-VERSION.N.msi` where `VERSION` is the MySQL Server version number such as 5.7 and `N` is the package number, which begins at 0.

MySQL Installer Commercial Edition

Download this edition from https://edelivery.oracle.com/ to install the commercial edition of all MySQL products for Windows. The Commercial Edition includes all of the products in the Community Edition and also includes the following products:

- Workbench SE/EE
- MySQL Enterprise Backup
- MySQL Enterprise Firewall

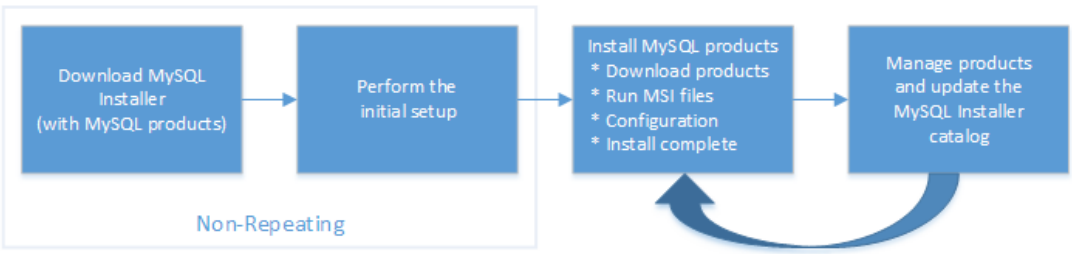
This edition integrates with your My Oracle Support (MOS) account. For knowledge-base content and patches, see My Oracle Support.

2.3.3.1 MySQL Installer Initial Setup

- MySQL Installer Licensing and Support Authentication
- Choosing a Setup Type
- Path Conflicts
- Check Requirements
- MySQL Installer Configuration Files

When you download MySQL Installer for the first time, a setup wizard guides you through the initial installation of MySQL products. As the following figure shows, the initial setup is a one-time activity in the overall process. MySQL Installer detects existing MySQL products installed on the host during its initial setup and adds them to the list of products to be managed.

Figure 2.7 MySQL Installer Process Overview



MySQL Installer extracts configuration files (described later) to the hard drive of the host during the initial setup. Although MySQL Installer is a 32-bit application, it can install both 32-bit and 64-bit binaries.

The initial setup adds a link to the Start menu under the **MySQL** group. Click **Start**, **All Programs**, **MySQL**, **MySQL Installer** to open MySQL Installer.

MySQL Installer Licensing and Support Authentication

MySQL Installer requires you to accept the license agreement before it will install new MySQL packages. After you accept the terms of the agreement, you can add, update, reconfigure, and remove all of the products and features provided by the MySQL Installer edition you downloaded.

For the commercial edition, entering your My Oracle Support (MOS) credentials is optional when installing bundled MySQL products, but your credentials are required when choosing unbundled MySQL products that MySQL Installer must download. An unbundled product is any MSI file that you download using MySQL Installer after the initial setup. Your credentials must match the user name and password that you have registered with Oracle for access to the support site.

Choosing a Setup Type

During the initial setup, you are prompted to select the MySQL products to be installed on the host. One alternative is to use a predetermined setup type that matches your setup requirements. Choosing one of the following setup types determines the initial installation only and does not limit your ability to install or update MySQL products for Windows later:

- **Developer Default**: Install the following products that compliment application development with MySQL:
  - MySQL Server (Installs the version that you selected when you downloaded MySQL Installer.)
  - MySQL Shell
  - MySQL Router
  - MySQL Workbench
  - MySQL for Visual Studio
  - MySQL for Excel
  - MySQL Notifier
  - MySQL Connectors (.NET / Python / ODBC / Java / C / C++)
  - MySQL Utilities
  - MySQL Documentation
  - MySQL Samples and Examples
- **Server only**: Only install the MySQL server. This setup type installs the general availability (GA) or development release server that you selected when you downloaded MySQL Installer. It uses the default installation and data paths.
- **Client only**: Only install the most recent MySQL applications and MySQL connectors. This setup type is similar to the `Developer Default` type, except that it does not include MySQL server or the client programs typically bundled with the server, such as `mysql` or `mysqladmin`.
- **Full**: Install all available MySQL products.
- **Custom** The custom setup type enables you to filter and select individual MySQL products from the MySQL Installer catalog.

Use the `Custom` setup type to install:

- A product or product version that is not available from the usual download locations. The catalog contains all product releases, including the other releases between pre-release (or development) and GA.
- An instance of MySQL server using an alternative installation path, data path, or both. For instructions on how to adjust the paths, see Setting Alternative Server Paths.
- Two or more MySQL server versions on the same host at the same time (for example, 5.6, 5.7, and 8.0).
- A specific combination of products and features not offered as a predetermine setup type. For example, you can install a single product, such as MySQL Workbench, instead of installing all client applications for Windows.

Path Conflicts

When the default installation or data folder (required by MySQL server) for a product to be installed already exists on the host, the wizard displays the **Path Conflict** step to identify each conflict and enable you to take action to avoid having files in the existing folder overwritten by the new installation. You see this step in the initial setup only when MySQL Installer detects a conflict.

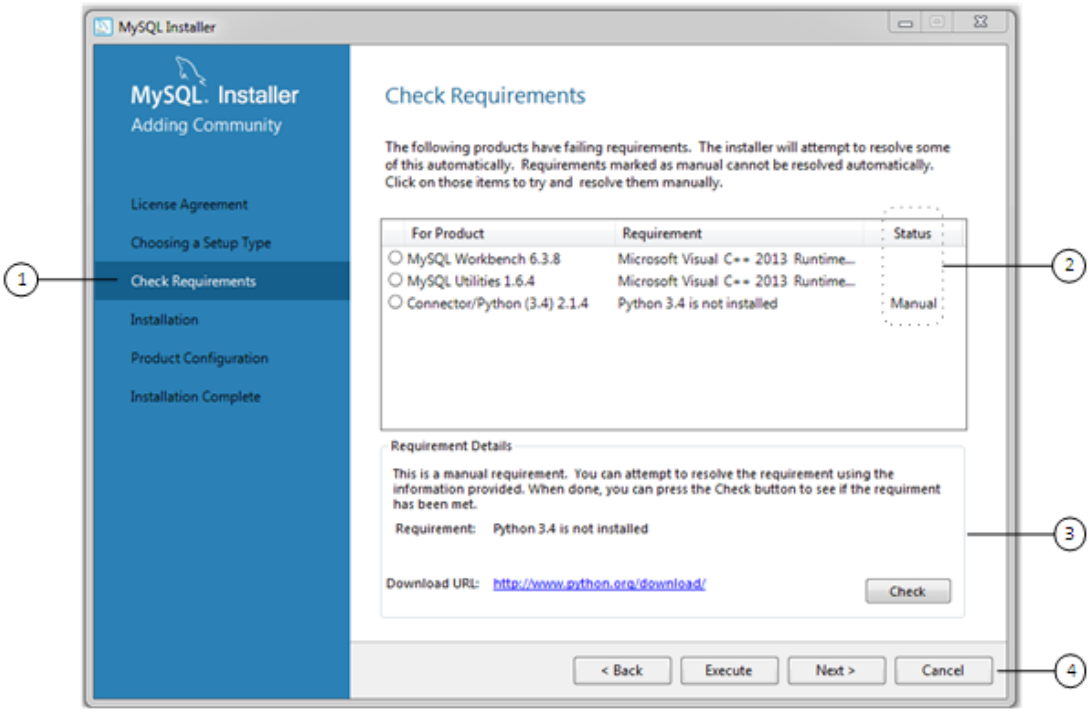
To resolve the path conflict, do one of the following:

- Select a product from the list to display the conflict options. A warning symbol indicates which path is in conflict. Use the browse button to choose a new path and then click **Next**.
- Click **Back** to choose a different setup type or product version, if applicable. The `Custom` setup type enables you to select individual product versions.
- Click **Next** to ignore the conflict and overwrite files in the existing folder.
- Delete the existing product. Click **Cancel** to stop the initial setup and close MySQL Installer. Open MySQL Installer again from the Start menu and delete the installed product from the host using the Delete operation from the dashboard.

Check Requirements

MySQL Installer uses entries in the `package-rules.xml` file to determine whether the prerequisite software for each product is installed on the host. When the requirements check fails, MySQL Installer displays the **Check Requirements** step to help you update the host. The following figure identifies and describes the key areas of this step.

Figure 2.8 Check Requirements



Description of Check Requirements Elements

1. Shows the current step in the initial setup. Steps in this list may change slightly depending on the products already installed on the host, the availability of prerequisite software, and the products to be installed on the host.

2. Lists all pending installation requirements by product and indicates the status as follows:

   - A blank space in the **Status** column means that MySQL Installer can attempt to download and install the required software for you.

   - The word *Manual* in the **Status** column means that you must satisfy the requirement manually. Select each product in the list to see its requirement details.

3. Describes the requirement in detail to assist you with each manual resolution. When possible, a download URL is provided. After you download and install the required software, click **Check** to verify that the requirement has been met.

4. Provides the following set operations to proceed:

   - **Back** – Return to the previous step. This action enables you to select a different the setup type.

   - **Execute** – Have MySQL Installer attempt to download and install the required software for all items without a manual status. Manual requirements are resolved by you and verified with the **Check** button.

   - **Next** – Do not execute the request to apply the requirements automatically and proceed to the installation without including the products that fail the check requirements step.

   - **Cancel** – Stop the installation of MySQL products. Because MySQL Installer is already installed, the initial setup begins again when you open MySQL Installer from the Start menu and click **Add** from the dashboard. For a description of the available management operations, see Product Catalog.

MySQL Installer Configuration Files

All MySQL Installer files are located within the `C:\Program Files (x86)` and `C:\ProgramData` folders. The following table describes the files and folders that define MySQL Installer as a standalone application.

Note

Installed MySQL products are neither altered nor removed when you update or uninstall MySQL Installer.

Table 2.5 MySQL Installer Configuration Files

File or Folder	Description	Folder Hierarchy
`MySQL Installer for Windows`	This folder contains all of the files needed to run MySQL Installer and MySQLInstallerConsole.exe, a command-line program with similar functionality.	`C:\Program Files (x86)`
`Templates`	The `Templates` folder has one file for each version of MySQL server. Template files contain keys and formulas to calculate some values dynamically.	`C:\ProgramData\MySQL \MySQL Installer for Windows\Manifest`
`package-rules.xml`	This file contains the prerequisites for every product to be installed.	`C:\ProgramData\MySQL \MySQL Installer for Windows\Manifest`
`produts.xml`	The `products` file (or product catalog) contains a list of all products available for download.	`C:\ProgramData\MySQL \MySQL Installer for Windows\Manifest`
`Product Cache`	The `Product Cache` folder contains all standalone MSI files bundled with the full package or downloaded afterward.	`C:\ProgramData\MySQL \MySQL Installer for Windows`

2.3.3.2 Installation Workflow

MySQL Installer provides a wizard-like tool to install and configure new MySQL products for Windows. Unlike the initial setup, which runs only once, MySQL Installer invokes the wizard each time you download or install a new product. For first-time installations, the steps of the initial setup proceed directly into the steps of the installation.

Note

Full permissions are granted to the user executing MySQL Installer to all generated files, such as `my.ini`. This does not apply to files and directories for specific products, such as the MySQL server data directory in `%ProgramData%` that is owned by `SYSTEM`.

Products installed and configured on a host follow a general pattern that might require your input during the various steps. MySQL Installer loads all selected products together using the following workflow:

- Product download. If you installed the full (not web) MySQL Installer package, all MSI files were loaded to the `Product Cache` folder during the initial setup and are not downloaded again. Otherwise, the status of each product changes from `Downloading` to `Downloaded`.
- Product installation. The status of each product in the list changes from `Ready to Install` to `Installing` to `Complete`. During the process, click **Show Details** to view the installation actions.

If you cancel the installation at this point, the products are installed, but the server (if installed) is not yet configured. To restart the server configuration, open MySQL Installer from the Start menu and click the **Reconfigure** link next to the appropriate server in the dashboard.

- Product configuration. This step applies to MySQL server and samples in most cases. The status for each item in the list should indicate, `Ready to Configure`. Click **Next** to begin the step-by-step configuration of all items in the list.

  The configuration options presented during this step depend on which version of the database you selected to install. After the installation completes, you can reconfigure MySQL server from the MySQL Installer dashboard.

- Installation compete. This step finalizes the installation and enables you to start some applications when the installation finishes.

InnoDB Cluster Sandbox Test Setup

You have two options to implement a high-availability solution when you install MySQL Server 5.7.17 or higher using MySQL Installer:

- Standalone MySQL Server / Classic MySQL Replication (default)

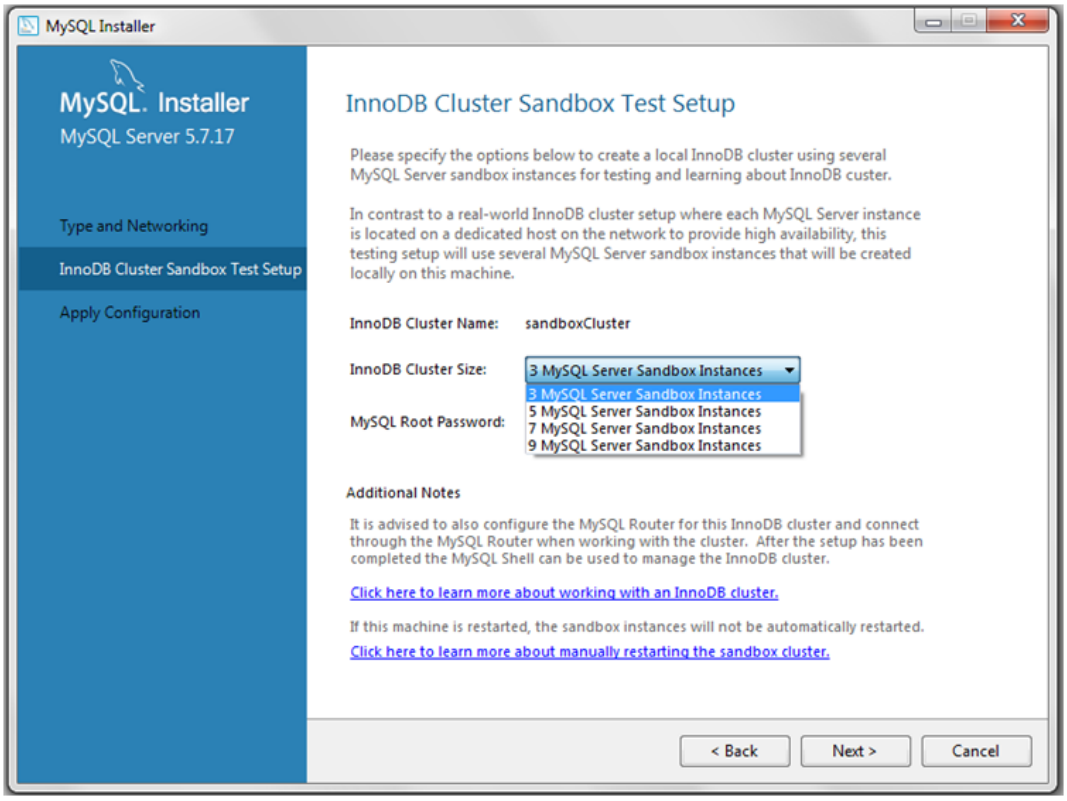
  Select this option to begin the initial configuration of a standalone MySQL server. You can configure multiple servers with classic MySQL Replication manually or use MySQL Shell 1.0.9 or higher to configure a production InnoDB cluster.

  For a description of the server configuration options that apply to a standalone MySQL server on Windows, see Server Configuration.

- InnoDB Cluster Sandbox Test Setup (for testing only)

  Select this option to create and configure InnoDB cluster sandbox instances locally for testing. You can configure an InnoDB cluster sandbox to have three, five, seven, or nine MySQL server instances. Use the **Reconfigure** quick action in the MySQL Installer toolbar to adjust the number of instances in the InnoDB cluster after the configuration has finished.

Figure 2.9 InnoDB Cluster Sandbox Test Setup



The InnoDB cluster sandbox, named `sandboxCluster` by default, is available on selected ports. After the configuration executes, click the **Summary** tab to view the specific ports that apply to your cluster.

You can use MySQL Installer to install MySQL Shell 1.0.9, if it is not installed. MySQL Shell enables you to manage the sandbox instances. To connect with the MySQL Shell on port 3310, execute the following command:

```
shell> mysqlsh root@localhost:3310
```

MySQL Installer also provides a wizard for configuring MySQL Router to connect to the test InnoDB cluster that was created in this step. For configuration details, see MySQL Router Configuration. To learn more about MySQL Router operations, see Routing for MySQL InnoDB cluster.

Server Configuration

MySQL Installer handles the initial configuration of the MySQL server. For example:

- It creates the configuration file (`my.ini`) that is used to configure the MySQL server. The values written to this file are influenced by choices you make during the installation process.

  Note

  Some definitions are host dependent. For example, query_cache is enabled if the host has fewer than three cores.

- By default, a Windows service for the MySQL server is added.
- Provides default installation and data paths for MySQL server. For instructions on how to change the default paths, see Setting Alternative Server Paths.

- It can optionally create MySQL server user accounts with configurable permissions based on general roles, such as DB Administrator, DB Designer, and Backup Admin. It optionally creates a Windows user named `MysqlSys` with limited privileges, which would then run the MySQL Server.

  User accounts may also be added and configured in MySQL Workbench.

- Checking **Show Advanced Options** allows additional **Logging Options** to be set. This includes defining custom file paths for the error log, general log, slow query log (including the configuration of seconds it requires to execute a query), and the binary log.

During the configuration process, click **Next** to proceed to the next step or **Back** to return to the previous step. Click **Execute** at the final step to apply the server configuration.

The sections that follow describe the server configuration options that apply to MySQL server on Windows. The server version you installed will determine which steps and options you can configure.

- Type and Networking
- Accounts and Roles
- Windows Service
- Plugins and Extensions
- Advanced Options
- Apply Server Configuration

Type and Networking

- Server Configuration Type

  Choose the MySQL server configuration type that describes your setup. This setting defines the amount of system resources (memory) that will be assigned to your MySQL server instance.

  - **Development**: A machine that will host many other applications, and typically this is your personal workstation. This option configures MySQL to use the least amount of memory.
  - **Server**: Several other applications will be running on this machine, such as a web server. This option configures MySQL to use a medium amount of memory.
  - **Dedicated**: A machine that is dedicated to running the MySQL server. Because no other major applications will run on this server, such as a web server, this option configures MySQL to use the majority of available memory.

- Connectivity

  Connectivity options control how the connection to MySQL is made. Options include:

  - **TCP/IP**: You may enable TCP/IP Networking here as otherwise only local host connections are allowed. Also define the **Port Number** and whether to open the firewall port for network access. If the port number is in use already, you will see the information icon ( ) next to the default value and **Next** is disabled until you provide a new port number.
  - **Named Pipe**: Enable and define the pipe name, similar to using the `--enable-named-pipe` option.
  - **Shared Memory**: Enable and then define the memory name, similar to using the `--shared-memory` option.

- Advanced Configuration

Check **Show Advanced Options** to set additional logging options in a later step. This includes defining custom file paths for the error log, general log, slow query log (including the configuration of seconds it requires to execute a query), and the binary log.

- MySQL Enterprise Firewall (Commercial Edition only)

  The **Enable Enterprise Firewall** check box is selected by default. For post-installation instructions, see MySQL Enterprise Firewall.

Accounts and Roles

- Root Account Password

  Assigning a root password is required and you will be asked for it when performing other MySQL Installer operations. Password strength is evaluated when you repeat the password in the box provided. For descriptive information regarding password requirements or status, move your mouse pointer over the information icon ( ) when it appears.

- MySQL User Accounts

  Optionally, you can create additional MySQL user accounts with predefined user roles. Each predefined role, such as DB Admin, are configured with their own set of privileges. For example, the DB Admin role has more privileges than the DB Designer role. Click the **Role** drop-down list for a description of each role.

Note

If the MySQL server is installed, then you must also enter the current root password.

Windows Service

On the Windows platform, MySQL server can run as a named service managed by the operating system and be configured to start up automatically when Windows starts. Alternatively, you can configure MySQL server to run as an executable program that requires manual configuration.

- **Configure MySQL server as a Windows service** (Selected by default.)

  When the default configuration option is selected, you can also select the following:

  - **Start the MySQL Server at System Startup**

    When selected (default), the service startup type is set to Automatic; otherwise, the startup type is set to Manual.

  - **Run Windows Service as**

    When **Standard System Account** is selected (default), the service logs on as Network Service.

    The **Custom User** option must have privileges to log on to Microsoft Windows as a service. The **Next** button will be disabled until this user is configured with the required privileges.

    A custom user is configured in Windows by searching for "local security policy" in the Start menu. In the Local Security Policy window, select **Local Policies**, **User Rights Assignment**, and then **Log On As A Service** to open the property dialog. Click **Add User or Group** to add the custom user and then click **OK** in each dialog to save the changes.

- Deselect the Windows Service option

Plugins and Extensions

The **Plugins and Extensions** step is visible during a new installation of MySQL server version 5.7.12 (or higher) only. If you are upgrading from a previous MySQL version, then you need to open MySQL Installer again and select the **Reconfigure** MySQL server option.

- **Enable X Protocol / MySQL as a Document Store** (Selected by default.)

  When the X Protocol option is selected, MySQL Installer loads and starts the X Plugin. Without the X Plugin running, X Protocol clients cannot connect to the server.

  - **Port Number:** *33060*

    Requires an unused port. The default port number is 33060.

  - **Open Firewall port for network access**

    Open by default when the X Protocol is selected.

  For more information about using MySQL as a document store and the X Plugin, see Section 20.2, “Key Concepts” and Section 20.7, “X Plugin”.

Advanced Options

This step is available if the **Show Advanced Configuration** check box was selected during the **Type and Networking** step. To enable this step now, click **Back** to return to the **Type and Networking** step and select the check box.

Advanced configuration options are related to the following MySQL log files:

- Error log
- General log
- Slow query log
- Bin log

Apply Server Configuration

All configuration settings are applied to the MySQL server when you click **Execute**. Use the **Configuration Steps** tab to follow the progress of each action; the icon for each toggles from white to green on success. Otherwise, the process stops and displays an error message if an individual action times out. Click the **Log** tab to view the log.

When the installation is done and you click **Finish**, MySQL Installer and the installed MySQL products are added to the Microsoft Windows Start menu under the `MySQL` group. Opening MySQL Installer loads the dashboard where installed MySQL products are listed and other MySQL Installer operations are available.

Setting Alternative Server Paths

You can change the default installation path, the data path, or both when you install MySQL server. After you have installed the server, the paths cannot be altered without removing and reinstalling the server instance.

To change paths for MySQL server

1. Identify the MySQL server to change and display the **Advanced Options** link.

   a. Navigate to the **Select Products and Features** step by doing one of the following:

i. If this is an initial setup, select the `Custom` setup type and click **Next**.

ii. If MySQL Installer is installed already, launch it from the Start menu and then click **Add** from the dashboard.

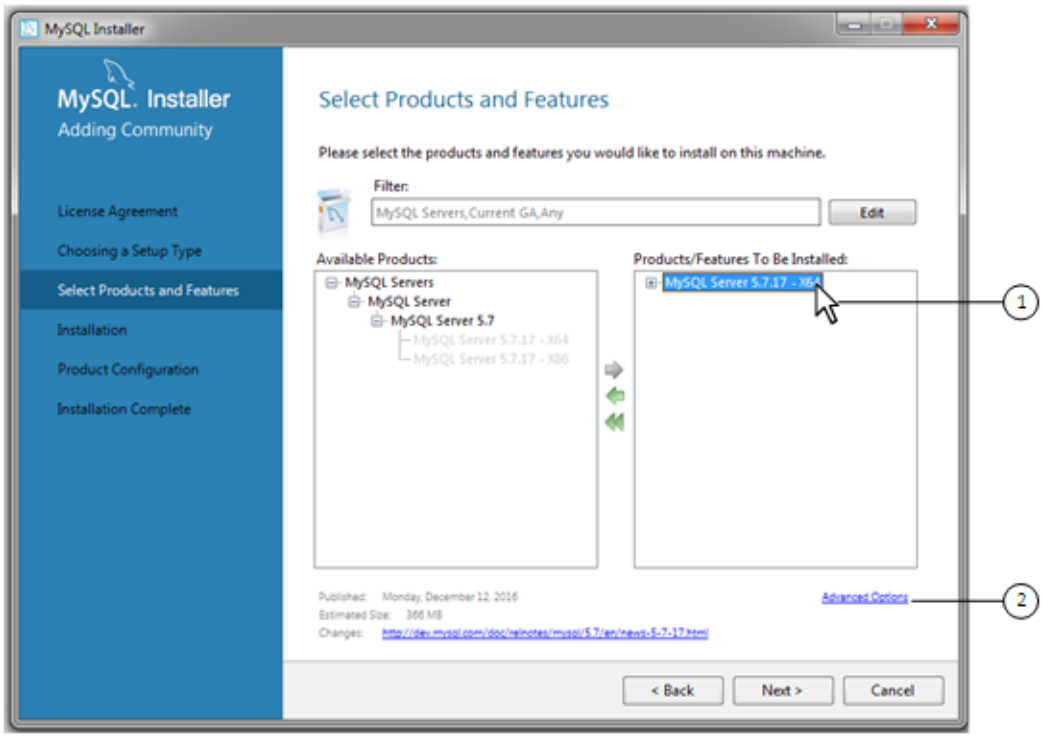
b. Click **Edit** to filter the list of products, locate the server instance to be installed in the **Available Products** list.

c. With the server instance selected, use the arrow to move the selected server to the **Products/Features To Be Installed** list.

d. Click the server to select it. When you select the server, the **Advanced Options** link appears. For details, see the figure that follows.

2. Click **Advanced Options** to open a dialog window with the path-setting options. After setting the path, click **Next** to continue with the configuration steps.

Figure 2.10 Change MySQL Server Path



MySQL Applications, Connectors, and Documentation

MySQL Installer provides you with a suite of tools for developing and managing business-critical applications on Windows. The suite consist of applications, connectors, documentation, and samples.

During the initial setup, choose any predetermined setup type, except `Server only`, to install the latest GA version of the tools. Use the `Custom` setup type to install an individual tool or specific version. If MySQL Installer is installed on the host already, use the **Add** operation to select and install tools from the MySQL Installer dashboard.

MySQL Router Configuration

MySQL Installer provides a configuration wizard that can bootstrap an installed instance of MySQL Router 2.1.3 or later to route traffic between MySQL applications and an InnoDB cluster. When configured, MySQL Router runs as a local Windows service. For detailed information about using MySQL Router with an InnoDB cluster, see Routing for MySQL InnoDB cluster.

To configure MySQL Router, do the following:

1. Set up InnoDB cluster. For instructions on how to configure an InnoDB cluster sandbox on the local host using MySQL Installer, see InnoDB Cluster Sandbox Test Setup. InnoDB cluster requires MySQL Server 5.7.17 or higher.

   For general InnoDB cluster information, see InnoDB Cluster.

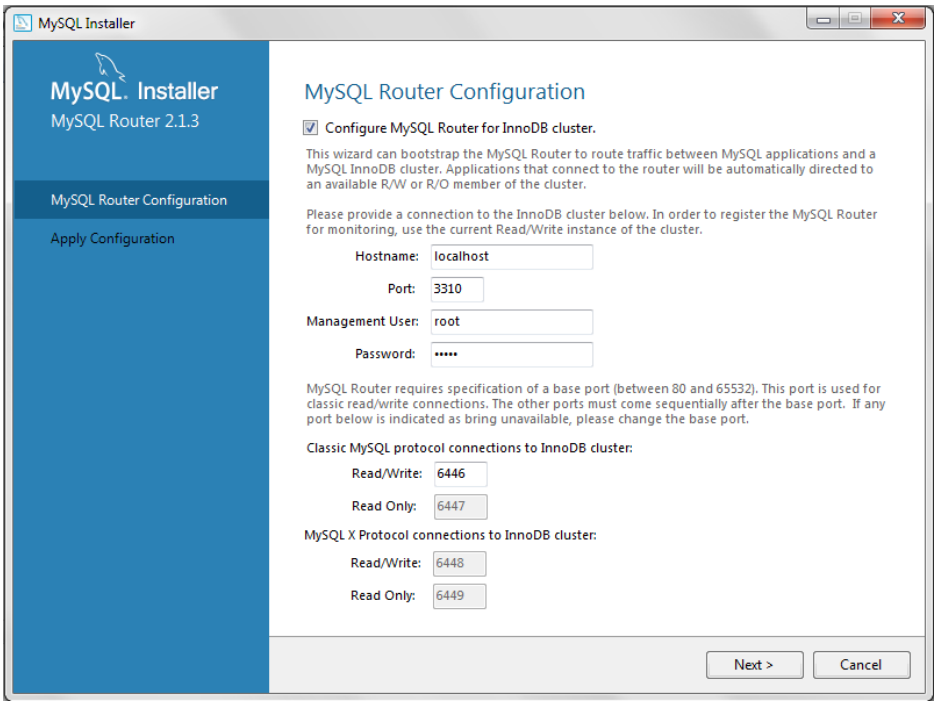
2. Using MySQL Installer, download and install the MySQL Router application. After the installation finishes, the configuration wizard prompts you for information. Select the **Configure MySQL Router for InnoDB cluster** check box to begin the configuration and provide the following configuration values:

   - **Hostname:** localhost by default.
   - **Port:** The port number of the primary server in the InnoDB cluster. The default is 3310.
   - **Management User:** An administrative user with root-level privileges.
   - **Password:** The password for the management user.
   - **Classic MySQL protocol connections to InnoDB cluster**

     Read/Write: Set the first base port number to one that is unused (between 80 and 65532) and the wizard will select the remaining ports for you.

     The figure that follows shows an example of the MySQL Router configuration screen.

   Figure 2.11 MySQL Router Configuration

3. Click **Next** and then **Execute** to apply the configuration. Click **Finish** to close MySQL Installer or return to the MySQL Installer dashboard.

2.3.3.3 MySQL Installer Product Catalog and Dashboard

- Product Catalog
- MySQL Installer Dashboard

- Locating Products to Install

This section describes the MySQL Installer product catalog and the dashboard.

Product Catalog

The product catalog stores the complete list of released MySQL products for Microsoft Windows that are available to download from MySQL. By default, and when an Internet connection is present, MySQL Installer updates the catalog daily. You can also update the catalog manually from the dashboard (described later).

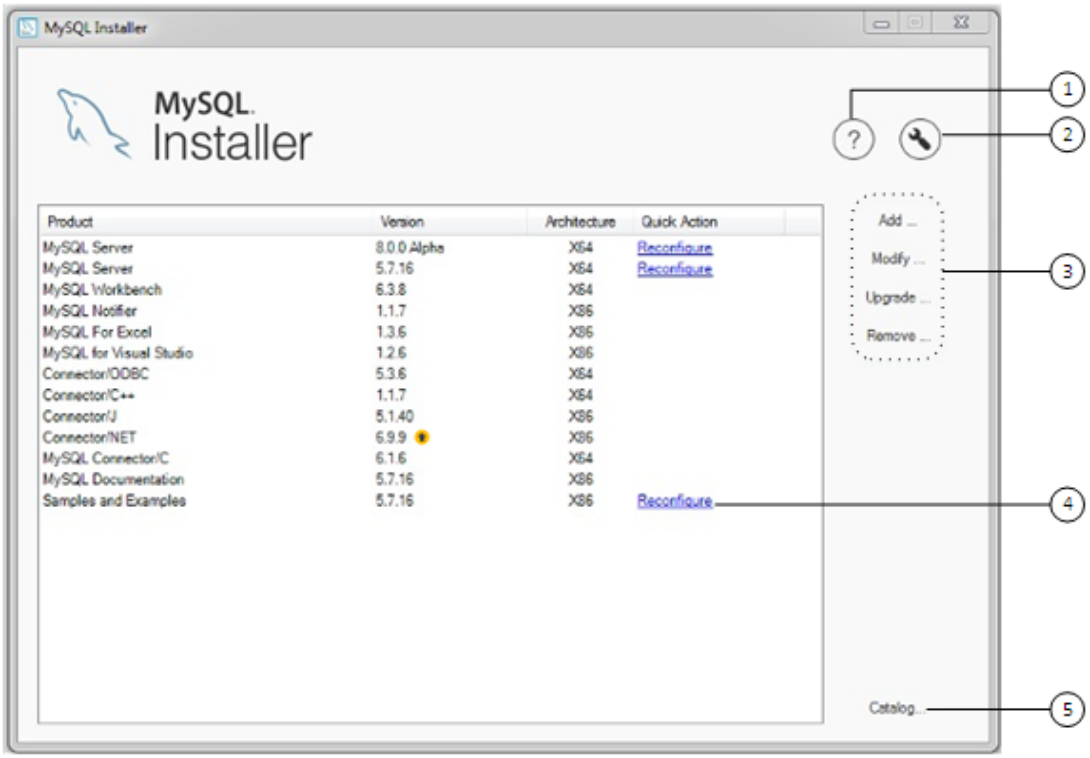
An up-to-date catalog performs the following actions:

- Populates the **Available Products** pane of the Select Products and Features step. This step appears when you select:
  - The `Custom` setup type during the initial setup.
  - The **Add** operation from the dashboard.
- Identifies when product updates are available for the installed products listed in the dashboard.

The catalog includes all development releases (Pre-Release), general releases (Current GA), and minor releases (Other Releases). Products in the catalog will vary somewhat, depending on the MySQL Installer edition that you download.

MySQL Installer Dashboard

Figure 2.12 MySQL Installer Dashboard Elements



Description of MySQL Installer Dashboard Elements

1. The About MySQL button (?) shows the current version of MySQL Installer and general information about MySQL. The version number is located above the **Back** button.

Always include this version number when reporting a problem with MySQL Installer.

2. The MySQL Installer Options button ( ) enables you to schedule daily automatic catalog updates. By default, catalog updates are scheduled at the hour when MySQL Installer was first installed. When new products or product versions are available, MySQL Installer adds them to the catalog and then displays an arrow icon ( ) next to the version number of installed products listed in the dashboard.

   Use this option to enable or disable automatic catalog updates and to reset the time of day when the MySQL Installer updates the catalog automatically. For specific settings, see the task named `ManifestUpdate` in the Windows Task Scheduler.

3. MySQL Installer dashboard operations provide a variety of actions that apply to installed products or products listed in the catalog. To initiate the following operations, first click the operation link and then select the product or products to manage:

   - **Add**: This operation opens the Select Products and Features screen. From there, you can filter the product in the product catalog, select one or more products to download (as needed), and begin the installation. For hints about using the filter, see Locating Products to Install.
   - **Modify**: Use this operation to add or remove the features associated with installed products. Features that you can modify vary in complexity by product. When the **Program Shortcut** check box is selected, the product appears in the Start menu under the `MySQL` group.
   - **Upgrade**: This operation loads the Select Products to Upgrade screen and populates it with all the upgrade candidates. An installed product can have more than one upgrade version and requires a current product catalog.

     To choose a new product version:

     a. Confirm that the check box next to product name in the **Upgradeable Products** pane has a check mark. Deselect the products that you do not intend to upgrade at this time.

     b. Click a product in the list to highlight it. This action populates the **Upgradeable Versions** pane with the details of each available version for the selected product: version number, published date, and a `Changes` link to open the release notes for that version.

     MySQL Installer upgrades all of the selected products in one action. Click **Show Details** to view the actions performed by MySQL Installer.

   - **Remove** This operation opens the Remove Products screen and populates it with the MySQL products installed on the host. Select the MySQL products you want to remove (uninstall) and then click **Execute** to begin the removal process.

     To select products to remove, do one of the following:

     - Select the check box for one or more products.
     - Select the **Product** check box to select all products.

4. The **Reconfigure** link in the Quick Action column next to each installed server loads the current configuration values for the server and then cycles through all configuration steps enabling you to change the options and values. On completion, MySQL Installer stops the server, applies the configuration changes, and restarts the server for you. For a description of each configuration option, see Server Configuration.

   Installed `Samples and Examples` associated with a specific MySQL server version can be also be reconfigured to apply feature-configuration changes, if any. You must provide credentials with root privileges to reconfigure these items.

5. The **Catalog** link enables you to download the latest catalog of MySQL products manually and then to integrate those product changes with MySQL Installer. The catalog-download action does not perform an upgrade of the products already installed on the host. Instead, it returns to the dashboard and displays an arrow icon in the Version column for each installed product that has a newer version. Use the **Upgrade** operation to install the newer product version.

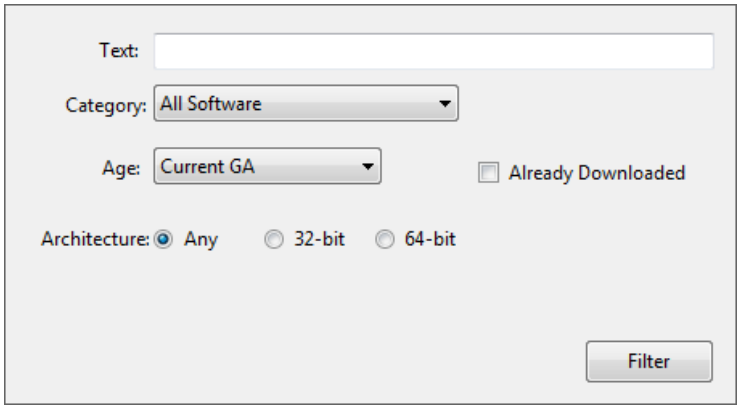
   You can also use the **Catalog** link to display the current change history of each product without downloading the new catalog. Select the **Do not update at this time** check box to view the change history only.

Locating Products to Install

MySQL products in the catalog are listed by category: MySQL Servers, Applications, MySQL Connectors, and Documentation. Only the latest GA versions appear in the **Available Products** pane by default. If you are looking for a pre-GA or older version of a product, it may not be visible in the default list.

To change the filter that displays available products in the Select Products and Features screen, click **Edit**.

Figure 2.13 Filter Available Products



Reset the following values to filter the list of available products:

- Text: Filter by text.
- Category: All Software (default), MySQL Servers, Applications, MySQL Connectors, or Documentation (for samples and documentation).
- Age: Pre-Release, Current GA (default), or Other Releases.
- Already Downloaded (the check box is deselected by default).
- Architecture: Any (default), 32-bit, or 64-bit.

2.3.3.4 MySQLInstallerConsole Reference

`MySQLInstallerConsole.exe` provides command-line functionality that is similar to MySQL Installer. It is installed when MySQL Installer is initially executed and then available within the `MySQL Installer` directory. Typically, that is in `C:\Program Files (x86)\MySQL\MySQL Installer \`, and the console must be executed with administrative privileges.

To use, invoke the command prompt with administrative privileges by choosing **Start**, **Accessories**, then right-click on **Command Prompt** and choose `Run as administrator`. And from the command line, optionally change the directory to where `MySQLInstallerConsole.exe` is located:

```
C:\> cd Program Files (x86)\MySQL\MySQL Installer for Windows
```

```
C:\Program Files (x86)\MySQL\MySQL Installer for Windows> MySQLInstallerConsole.exe help
=================== Start Initialization ===================
MySQL Installer is running in Community mode

Attempting to update manifest.
Initializing product requirements
Loading product catalog
Checking for product catalog snippets
Checking for product packages in the bundle
Categorizing product catalog
Finding all installed packages.
Your product catalog was last updated at 11/1/2016 4:10:38 PM
=================== End Initialization ===================

The following commands are available:

Configure - Configures one or more of your installed programs.
Help      - Provides list of available commands.
Install   - Install and configure one or more available MySQL programs.
List      - Provides an interactive way to list all products available.
Modify    - Modifies the features of installed products.
Remove    - Removes one or more products from your system.
Status    - Shows the status of all installed products.
Update    - Update the current product catalog.
Upgrade   - Upgrades one or more of your installed programs.
```

`MySQLInstallerConsole.exe` supports the following options, which are specified on the command line:

Note

Configuration block values that contain a colon (":") must be wrapped in double quotes. For example, installdir="C:\MySQL\MySQL Server 8.0".

- `configure` *`[product1]:[setting]=[value]; [product2]:[setting]=[value]; [...]`*

  Configure one or more MySQL products on your system. Multiple setting=value pairs can be configured for each product.

  Switches include:

  - `-showsettings` : Displays the available options for the selected product, by passing in the product name after `-showsettings`.
  - `-silent` : Disable confirmation prompts.

  ```
  C:\> MySQLInstallerConsole configure -showsettings server
  C:\> MySQLInstallerConsole configure server:port=3307
  ```

- `help` *`[command]`*

  Displays a help message with usage examples, and then exits. Pass in an additional command to receive help specific to that command.

  ```
  C:\> MySQLInstallerConsole help
  C:\> MySQLInstallerConsole help install
  ```

- `install` *`[product]:[features]:[config block]:[config block]:[config block]; [...]`*

  Install one or more MySQL products on your system.

  Switches and syntax options include:

- `-type=[SetupType]` : Installs a predefined set of software. The "SetupType" can be one of the following:

  Note

  Non-custom setup types can only be chosen if no other MySQL products are installed.

  - **Developer**: Installs a complete development environment.
  - **Server**: Installs a single MySQL server
  - **Client**: Installs client programs and libraries
  - **Full**: Installs everything
  - **Custom**: Installs user selected products. This is the default option.

- `-showsettings` : Displays the available options for the selected product, by passing in the product name after `-showsettings`.
- `-silent` : Disable confirmation prompts.
- `[config block]`: One or more configuration blocks can be specified. Each configuration block is a semicolon separated list of key value pairs. A block can include either a "config" or "user" type key, where "config" is the default type if one is not defined.

  Configuration block values that contain a colon (":") must be wrapped in double quotes. For example, `installdir="C:\MySQL\MySQL Server 8.0"`.

  Only one "config" type block can be defined per product. A "user" block should be defined for each user that should be created during the product's installation.

  Note

  Adding users is not supported when a product is being reconfigured.

- `[feature]`: The feature block is a semicolon separated list of features, or '*' to select all features.

```
C:\> MySQLInstallerConsole install server;5.6.25:*:port=3307;serverid=2:type=user;username=foo;passwo
C:\> MySQLInstallerConsole install server;5.6.25;x64 -silent
```

An example that passes in additional configuration blocks, broken up by `^` to fit this screen:

```
C:\> MySQLInstallerConsole install server;5.6.25;x64:*:type=config;openfirewall=true; ^
         generallog=true;binlog=true;serverid=3306;enable_tcpip=true;port=3306;rootpasswd=pass; ^
         installdir="C:\MySQL\MySQL Server 5.6":type=user;datadir="C:\MySQL\data";username=foo;passw
```

- `list`

  Lists an interactive console where all of the available MySQL products can be searched. Execute `MySQLInstallerConsole list` to launch the console, and enter in a substring to search.

```
C:\> MySQLInstallerConsole list
```

- `modify [product1:-removelist|+addlist] [product2:-removelist|+addlist] [...]`

  Modifies or displays features of a previously installed MySQL product.

- `-silent` : Disable confirmation prompts.

```
C:\> MySQLInstallerConsole modify server
C:\> MySQLInstallerConsole modify server:+documentation
C:\> MySQLInstallerConsole modify server:-debug
```

- `remove [product1] [product2] [...]`

  Removes one ore more products from your system.

  - `*` : Pass in `*` to remove all of the MySQL products.
  - `-continue` : Continue the operation even if an error occurs.
  - `-silent` : Disable confirmation prompts.

```
C:\> MySQLInstallerConsole remove *
C:\> MySQLInstallerConsole remove server
```

- `status`

  Provides a quick overview of the MySQL products that are installed on the system. Information includes product name and version, architecture, date installed, and install location.

```
C:\> MySQLInstallerConsole status
```

- `update`

  Downloads the latest MySQL product catalog to your system. On success, the download catalog will be applied the next time either MySQLInstaller or MySQLInstallerConsole is executed.

```
C:\> MySQLInstallerConsole update
```

Note

The **Automatic Catalog Update** GUI option executes this command from the Windows Task Scheduler.

- `upgrade [product1:version] [product2:version] [...]`

  Upgrades one or more products on your system. Syntax options include:

  - `*` : Pass in `*` to upgrade all products to the latest version, or pass in specific products.
  - `!` : Pass in `!` as a version number to upgrade the MySQL product to its latest version.
  - `-silent` : Disable confirmation prompts.

```
C:\> MySQLInstallerConsole upgrade *
C:\> MySQLInstallerConsole upgrade workbench:6.3.5
C:\> MySQLInstallerConsole upgrade workbench:!
C:\> MySQLInstallerConsole upgrade workbench:6.3.5 excel:1.3.2
```

2.3.4 MySQL Notifier

MySQL Notifier is a tool that enables you to monitor and adjust the status of your local and remote MySQL server instances through an indicator that resides in the system tray. MySQL Notifier also gives quick access to MySQL Workbench through its context menu.

The MySQL Notifier is installed by MySQL Installer, and (by default) will start-up when Microsoft Windows is started.

To install, download and execute the MySQL Installer, be sure the MySQL Notifier product is selected, then proceed with the installation. See the MySQL Installer manual for additional details.

For notes detailing the changes in each release of MySQL Notifier, see the MySQL Notifier Release Notes.

Visit the MySQL Notifier forum for additional MySQL Notifier help and support.

Features include:

- Start, Stop, and Restart instances of the MySQL Server.
- Automatically detects (and adds) new MySQL Server services. These are listed under **Manage Monitored Items**, and may also be configured.
- The Tray icon changes, depending on the status. It's green if all monitored MySQL Server instances are running, or red if at least one service is stopped. The **Update MySQL Notifier tray icon based on service status** option, which dictates this behavior, is enabled by default for each service.
- Links to other applications like MySQL Workbench, MySQL Installer, and the MySQL Utilities. For example, choosing **Manage Instance** will load the MySQL Workbench Server Administration window for that particular instance.
- If MySQL Workbench is also installed, then the **Manage Instance** and **SQL Editor** options are available for local (but not remote) MySQL instances.
- Monitors both local and remote MySQL instances.

2.3.4.1 MySQL Notifier Usage

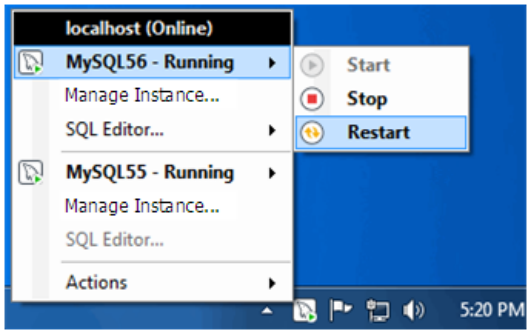
MySQL Notifier resides in the system tray and provides visual status information for your MySQL server instances. A green icon is displayed at the top left corner of the tray icon if the current MySQL server is running, or a red icon if the service is stopped.

MySQL Notifier automatically adds discovered MySQL services on the local machine, and each service is saved and configurable. By default, the **Automatically add new services whose name contains** option is enabled and set to `mysql`. Related **Notifications Options** include being notified when new services are either discovered or experience status changes, and are also enabled by default. And uninstalling a service will also remove the service from MySQL Notifier.

Clicking the system tray icon will reveal several options, as the follow figures show:

The Service Instance menu is the main MySQL Notifier window, and enables you to Stop, Start, and Restart the MySQL server.

Figure 2.14 MySQL Notifier Service Instance menu

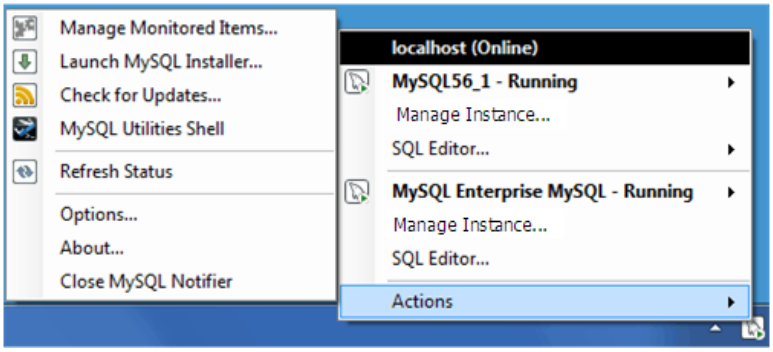


The **Actions** menu includes several links to external applications (if they are installed), and a **Refresh Status** option to manually refresh the status of all monitored services (in both local and remote computers) and MySQL instances.

Note

The main menu will not show the **Actions** menu when there are no services being monitored by MySQL Notifier.

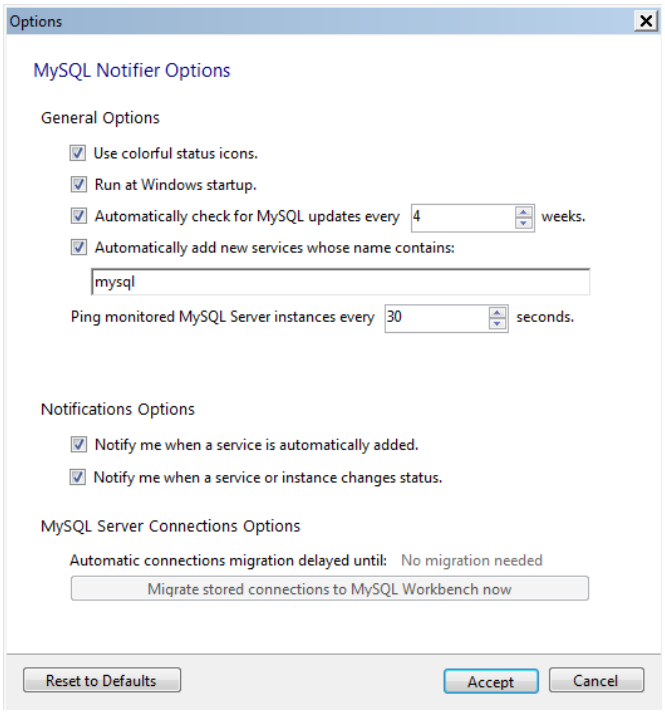
Figure 2.15 MySQL Notifier Actions menu



The **Actions**, **Options** menu configures MySQL Notifier and includes options to:

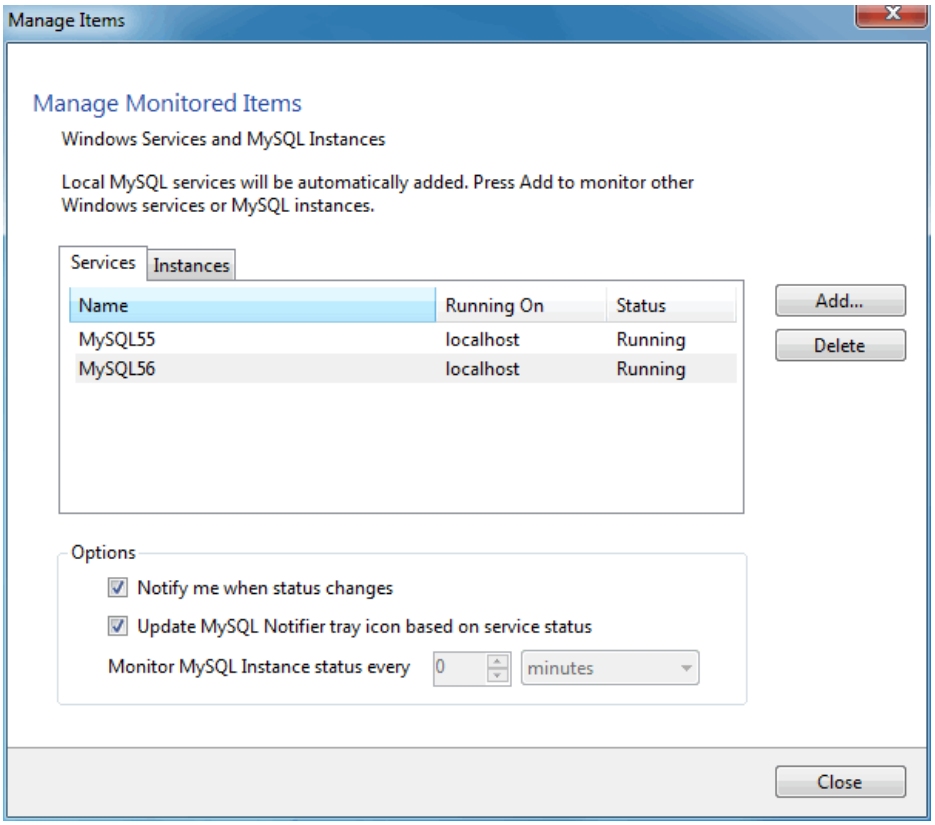
- **Use colorful status icons**: Enables a colorful style of icons for the tray of MySQL Notifier.
- **Run at Windows Startup**: Allows the application to be loaded when Microsoft Windows starts.
- **Automatically Check For Updates Every** # **Weeks**: Checks for a new version of MySQL Notifier, and runs this check every # weeks.
- **Automatically add new services whose name contains:** The text used to filter services and add them automatically to the monitored list of the local computer running MySQL Notifier, and on remote computers already monitoring Windows services.
- **Ping monitored MySQL Server instances every** # **seconds**: The interval (in seconds) to ping monitored MySQL Server instances for status changes. Longer intervals might be necessary if the list of monitored remote instances is large.
- **Notify me when a service is automatically added**: Will display a balloon notification from the taskbar when a newly discovered service is added to the monitored services list.
- **Notify me when a service changes status**: Will display a balloon notification from the taskbar when a monitored service changes its status.
- **Automatic connections migration delayed until:** When there are connections to migrate, postpone the migration by one hour, one day, one week, one month, or indefinitely.

Figure 2.16 MySQL Notifier Options menu



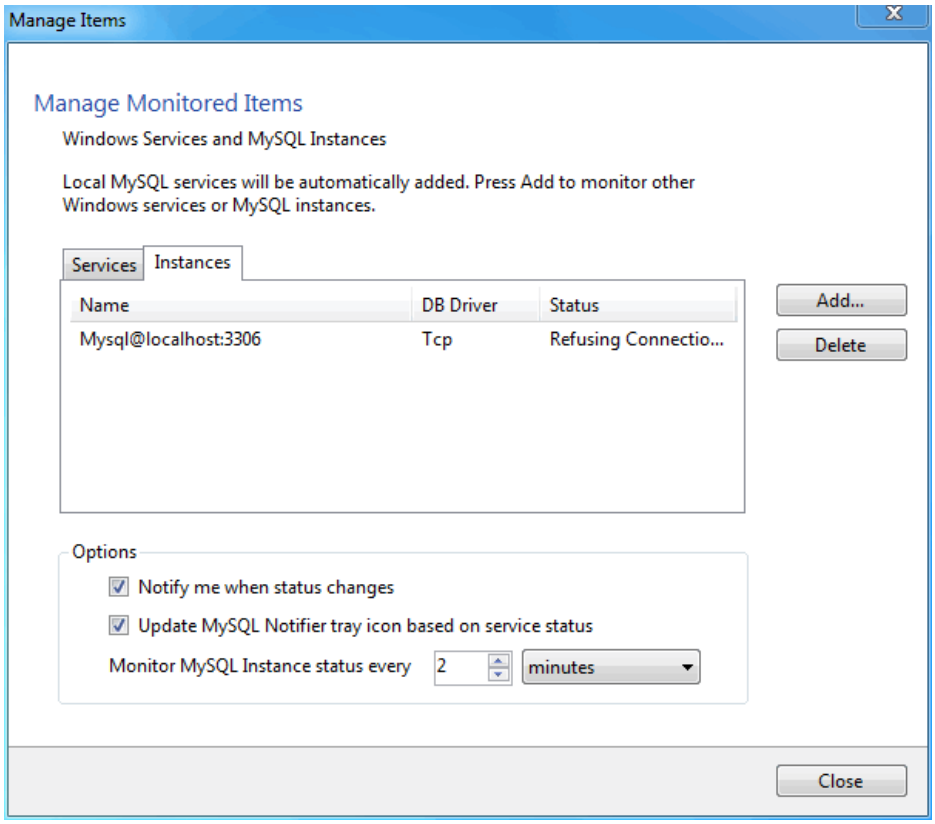
The **Actions**, **Manage Monitored Items** menu enables you to configure the monitored services and MySQL instances. First, with the **Services** tab open:

Figure 2.17 MySQL Notifier Manage Services menu



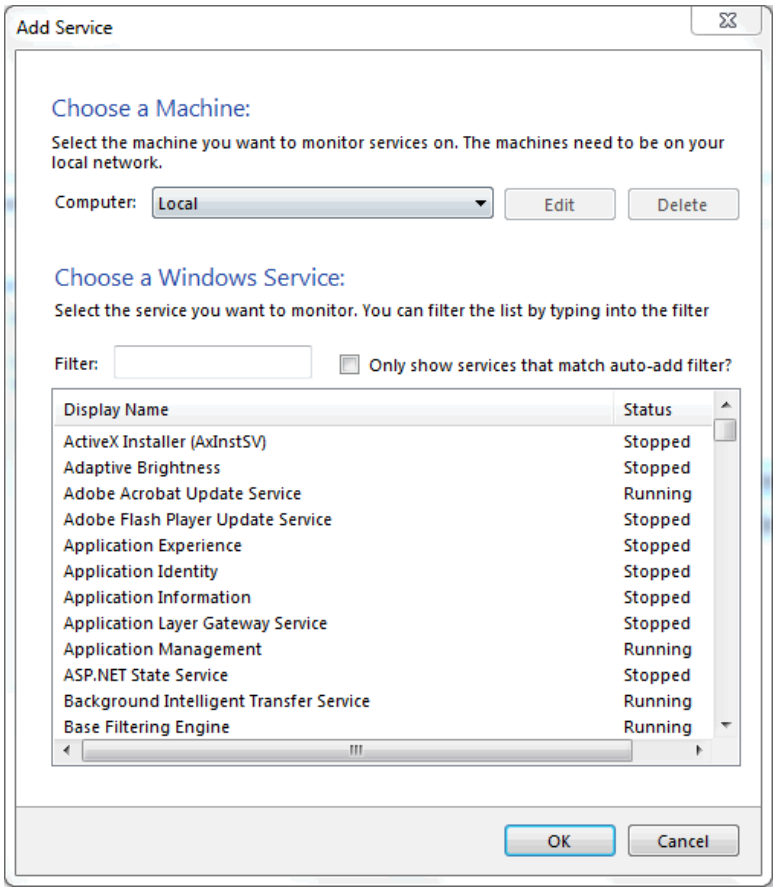
The **Instances** tab is similar:

Figure 2.18 MySQL Notifier Manage Instances menu



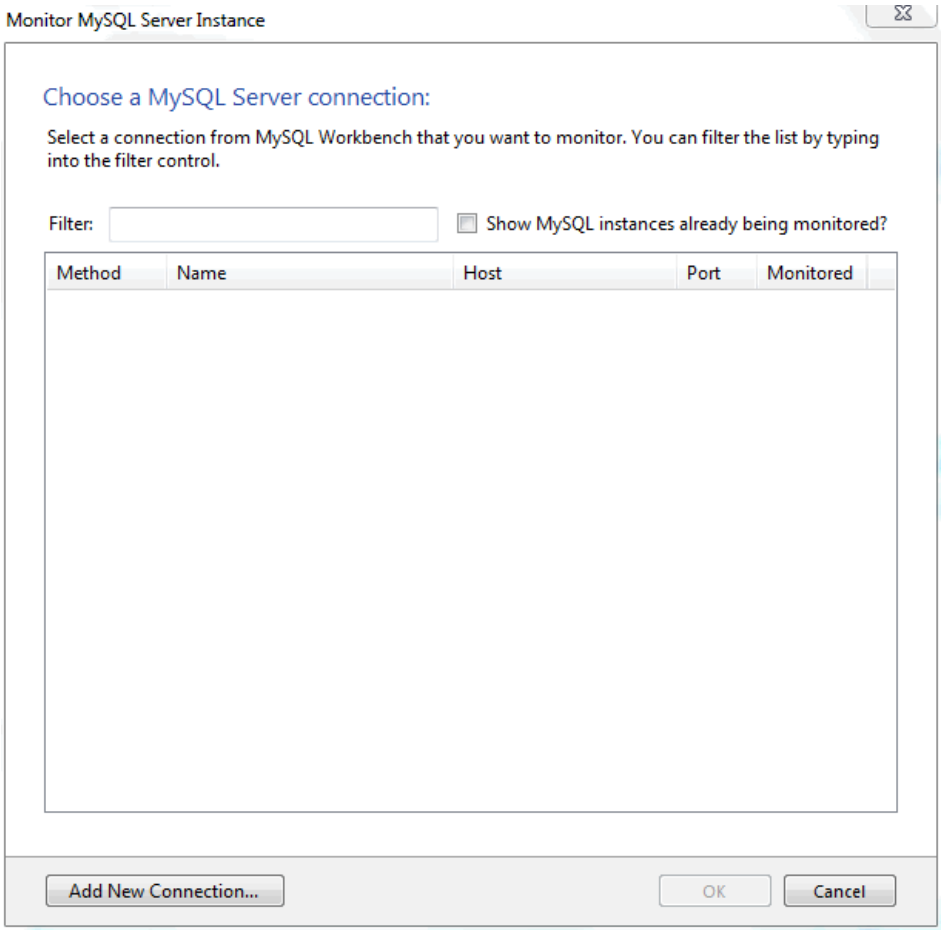
Adding a service or instance (after clicking **Add** in the **Manage Monitored Items** window) enables you to select a running Microsoft Windows service or instance connection, and configure MySQL Notifier to monitor it. Add a new service or instance by clicking service name from the list, then **OK** to accept. Multiple services and instances may be selected.

Figure 2.19 MySQL Notifier Adding new services



Add instances:

Figure 2.20 MySQL Notifier Adding new instances



Troubleshooting

For issues that are not documented here, visit the MySQL Notifier Support Forum for MySQL Notifier help and support.

- *Problem*: attempting to start/stop/restart a MySQL service might generate an error similar to "The Service **MySQL*VERSION*** failed the most recent status change request with the message "The service **mysql*VERSION*** was not found in the Windows Services".

  Explanation: this is a case-sensitivity issue, in that the service name is **MySQL*VERSION*** compared to having **mysql*VERSION*** in the configuration file.

  Solution: either update your MySQL Notifier configuration file with the correct information, or stop MySQL Notifier and delete this configuration file. The MySQL Notifier configuration file is located at `%APPDATA%\Oracle\MySQL Notifier\settings.config` where `%APPDATA%` is a variable and depends on your system. A typical location is "C:\Users*YourUsername*\AppData\Running\Oracle\MySQL Notifier\settings.config" where *YourUsername* is your system's user name. In this file, and within the ServerList section, change the ServerName values from lowercase to the actual service names. For example, change **mysql*VERSION*** to **MySQL*VERSION***, save, and then restart MySQL Notifier. Alternatively, stop MySQL Notifier, delete this file, then restart MySQL Notifier.

- *Problem*: when connecting to a remote computer for the purpose of monitoring a remote Windows service, the **Add Service** dialog does not always show all the services shown in the Windows Services console.

  Explanation: this behavior is governed by the operating system and the outcome is expected when working with nondomain user accounts. For a complete description of the behavior, see the User Account Control and WMI article from Microsoft.

Solution: when the remote computer is in a compatible domain, it is recommended that domain user accounts are used to connect through WMI to a remote computer. For detailed setup instructions using WMI, see Section 2.3.4.2, “Setting Up Remote Monitoring in MySQL Notifier”.

Alternatively, when domain user accounts are not available, Microsoft provides a less secure workaround that should only be implemented with caution. For more information, see the Description of User Account Control and remote restrictions in Windows Vista KB article from Microsoft.

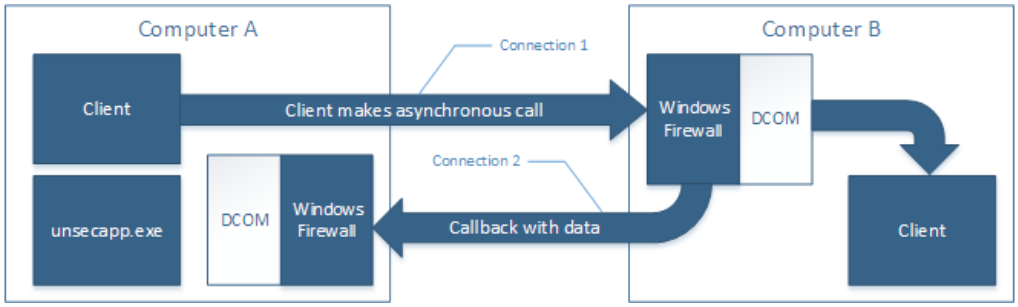
2.3.4.2 Setting Up Remote Monitoring in MySQL Notifier

MySQL Notifier uses Windows Management Instrumentation (WMI) to manage and monitor services on remote computers. This section explains how it works and how to set up your system to monitor remote MySQL instances.

In order to configure WMI, it is important to understand that the underlying Distributed Component Object Model (DCOM) architecture is doing the WMI work. Specifically, MySQL Notifier is using asynchronous notification queries on remote Microsoft Windows hosts as .NET events. These events send an asynchronous callback to the computer running MySQL Notifier so it knows when a service status has changed on the remote computer. Asynchronous notifications offer the best performance compared to semisynchronous notifications or synchronous notifications that use timers.

Asynchronous notification requires the remote computer to send a callback to the client computer (thus opening a reverse connection), so the Windows Firewall and DCOM settings must be properly configured for the communication to function properly.

Figure 2.21 MySQL Notifier Distributed Component Object Model (DCOM)



Most of the common errors thrown by asynchronous WMI notifications are related to Windows Firewall blocking the communication, or to DCOM / WMI settings not being set up properly. For a list of common errors with solutions, see Common Errors.

The following steps are required to make WMI function. These steps are divided between two machines. A single host computer that runs MySQL Notifier (Computer A), and multiple remote machines that are being monitored (Computer B).

Computer running MySQL Notifier (Computer A)

1. Enable remote administration by either editing the **Group Policy Editor**, or using `NETSH`:

   Using the **Group Policy Editor**:

   a. Click **Start**, click **Run**, type `GPEDIT.MSC`, and then click **OK**.

   b. Under the **Local Computer Policy** heading, expand **Computer Configuration**.

   c. Expand **Administrative Templates**, then **Network**, **Network Connections**, and then **Windows Firewall**.

   d. If the computer is in the domain, then double-click **Domain Profile**; otherwise, double-click **Standard Profile**.

e. Double-click **Windows Firewall: Allow inbound remote administration exception** to open a configuration window.

f. Check the **Enabled** option button and then click **OK**.

Using the `NETSH` command:

Note

The "netsh firewall" command is deprecated as of Microsoft Server 2008 and Vista, and replaced with "netsh advfirewall firewall".

a. Open a command prompt window with Administrative rights (you can right-click the Command Prompt icon and select **Run as Administrator**).

b. Execute the following command:

```
NETSH advfirewall firewall set service RemoteAdmin enable
```

2. Open the DCOM port TCP 135:

a. Open a command prompt window with Administrative rights (you can right-click the Command Prompt icon and select **Run as Administrator**).

b. Execute the following command:

```
NETSH advfirewall firewall add rule name=DCOM_TCP135 protocol=TCP localport=135 dir=in action=allow
```

3. Add the client application that contains the sink for the callback (`MySqlNotifier.exe`) to the Windows Firewall Exceptions List (use either the Windows Firewall configuration or `NETSH`):

Using the Windows Firewall configuration:

a. In the Control Panel, double-click **Windows Firewall**.

b. In the Windows Firewall window's left panel, click **Allow a program or feature through Windows Firewall**.

c. In the Allowed Programs window, click **Change Settings** and do one of the following:

- If `MySqlNotifier.exe` is in the Allowed programs and features list, make sure it is checked for the type of networks the computer connects to (Private, Public or both).
- If `MySqlNotifier.exe` is not in the list, click **Allow another program...**.

  i. In the **Add a Program** window, select the `MySqlNotifier.exe` if it exists in the Programs list, otherwise click Browse... and go to the directory where `MySqlNotifier.exe` was installed to select it, then click **Add**.

  ii. Make sure `MySqlNotifier.exe` is checked for the type of networks the computer connects to (Private, Public or both).

Using the `NETSH` command:

a. Open a command prompt window with Administrative rights (you can right-click the Command Prompt icon and click **Run as Administrator**).

b. Execute the following command, where you change "*`[YOUR_INSTALL_DIRECTORY]`*":

```
```

```
NETSH advfirewall firewall add rule name=MySqlNotifier program=[YOUR_INSTALL_DIRECTORY]\MySqlNot
```

4. If Computer B is either a member of `WORKGROUP` or is in a different domain that is untrusted by Computer A, then the callback connection (Connection 2) is created as an Anonymous connection. To grant Anonymous connections DCOM Remote Access permissions:

   a. Click **Start**, click **Run**, type `DCOMCNFG`, and then click **OK**.

   b. In the Component Services dialog box, expand Component Services, expand Computers, and then right-click **My Computer** and click **Properties**.

   c. In the My Computer Properties dialog box, click the **COM Security** tab.

   d. Under Access Permissions, click **Edit Limits**.

   e. In the Access Permission dialog box, select **ANONYMOUS LOGON name** in the Group or user names box. In the Allow column under Permissions for User, select **Remote Access**, and then click **OK**.

Monitored Remote Computer (Computer B)

If the user account that is logged on to the computer running the MySQL Notifier (Computer A) is a local administrator on the remote computer (Computer B), such that the same account is an administrator on Computer B, you can skip to the "Allow for remote administration" step.

Setting DCOM security to allow a non-administrator user to access a computer remotely:

1. Grant "DCOM remote launch" and activation permissions for a user or group:

   a. Click **Start**, click **Run**, type `DCOMCNFG`, and then click **OK**.

   b. In the Component Services dialog box, expand Component Services, expand Computers, and then right-click **My Computer** and click **Properties**.

   c. In the My Computer Properties dialog box, click the **COM Security** tab.

   d. Under Launch and Activation Permission, click **Edit Limits**.

   e. In the **Launch and Activation Permission** dialog box, follow these steps if your name or your group does not appear in the Groups or user names list:

      i. In the **Launch and Activation Permission** dialog box, click **Add**.

      ii. In the Select Users or Groups dialog box, add your name and the group in the **Enter the object names to select** box, and then click **OK**.

   f. In the **Launch and Activation Permission** dialog box, select your user and group in the Group or user names box. In the Allow column under Permissions for User, select **Remote Launch**, select **Remote Activation**, and then click **OK**.

   Grant DCOM remote access permissions:

   a. Click **Start**, click **Run**, type `DCOMCNFG`, and then click **OK**.

   b. In the Component Services dialog box, expand Component Services, expand Computers, and then right-click **My Computer** and click **Properties**.

   c. In the My Computer Properties dialog box, click the **COM Security** tab.

   d. Under Access Permissions, click **Edit Limits**.

e. In the Access Permission dialog box, select **ANONYMOUS LOGON name** in the Group or user names box. In the Allow column under Permissions for User, select **Remote Access**, and then click **OK**.

2. Allowing non-administrator users access to a specific WMI namespace:

   a. In the Control Panel, double-click **Administrative Tools**.

   b. In the Administrative Tools window, double-click **Computer Management**.

   c. In the Computer Management window, expand the **Services and Applications** tree.

   d. Right-click the WMI Control icon and select **Properties**.

   e. In the WMI Control Properties window, click the **Security** tab.

   f. In the Security tab, select the namespace and click **Security**. Root/CIMV2 is a commonly used namespace.

   g. Locate the appropriate account and check **Remote Enable** in the Permissions list.

3. Allow for remote administration by either editing the **Group Policy Editor** or using `NETSH`:

   Using the **Group Policy Editor**:

   a. Click **Start**, click **Run**, type `GPEDIT.MSC`, and then click **OK**.

   b. Under the Local Computer Policy heading, double-click **Computer Configuration**.

   c. Double-click **Administrative Templates**, then **Network**, **Network Connections**, and then **Windows Firewall**.

   d. If the computer is in the domain, then double-click **Domain Profile**; otherwise, double-click **Standard Profile**.

   e. Click **Windows Firewall: Allow inbound remote administration exception**.

   f. On the Action menu either select **Edit**, or double-click the selection from the previous step.

   g. Check the **Enabled** radio button, and then click **OK**.

   Using the `NETSH` command:

   a. Open a command prompt window with Administrative rights (you can right-click the Command Prompt icon and click Run as Administrator).

   b. Execute the following command:

   ```
   NETSH advfirewall firewall set service RemoteAdmin enable
   ```

4. Confirm that the user account you are logging in with uses the `Name` value and not the `Full Name` value:

   a. In the **Control Panel**, double-click **Administrative Tools**.

   b. In the **Administrative Tools** window, double-click **Computer Management**.

   c. In the **Computer Management** window, expand the **System Tools** then **Local Users and Groups**.

d. Click the **Users** node, and on the right side panel locate your user and make sure it uses the **Name** value to connect, and not the **Full Name** value.

Common Errors

- `0x80070005`
  - DCOM Security was not configured properly (see Computer B, the `Setting DCOM security...` step).
  - The remote computer (Computer B) is a member of WORKGROUP or is in a domain that is untrusted by the client computer (Computer A) (see Computer A, the `Grant Anonymous connections DCOM Remote Access permissions` step).
- `0x8007000E`
  - The remote computer (Computer B) is a member of WORKGROUP or is in a domain that is untrusted by the client computer (Computer A) (see Computer A, the `Grant Anonymous connections DCOM Remote Access permissions` step).
- `0x80041003`
  - Access to the remote WMI namespace was not configured properly (see Computer B, the `Allowing non-administrator users access to a specific WMI namespace` step).
- `0x800706BA`
  - The DCOM port is not open on the client computers (Computer A) firewall. See the `Open the DCOM port TCP 135` step for Computer A.
  - The remote computer (Computer B) is inaccessible because its network location is set to Public. Make sure you can access it through the Windows Explorer.

2.3.5 Installing MySQL on Microsoft Windows Using a noinstall Zip Archive

Users who are installing from the `noinstall` package can use the instructions in this section to manually install MySQL. The process for installing MySQL from a Zip archive is as follows:

1. Extract the main archive to the desired install directory

   Optional: also extract the debug-test archive if you plan to execute the MySQL benchmark and test suite

2. Create an option file
3. Choose a MySQL server type
4. Initialize MySQL
5. Start the MySQL server
6. Secure the default user accounts

This process is described in the sections that follow.

2.3.5.1 Extracting the Install Archive

To install MySQL manually, do the following:

1. If you are upgrading from a previous version please refer to Section 2.3.8, “Upgrading MySQL on Windows”, before beginning the upgrade process.

2. Make sure that you are logged in as a user with administrator privileges.

3. Choose an installation location. Traditionally, the MySQL server is installed in `C:\mysql`. The MySQL Installation Wizard installs MySQL under `C:\Program Files\MySQL`. If you do not install MySQL at `C:\mysql`, you must specify the path to the install directory during startup or in an option file. See Section 2.3.5.2, “Creating an Option File”.

Note

The MySQL Installer installs MySQL under `C:\Program Files\MySQL`.

4. Extract the install archive to the chosen installation location using your preferred Zip archive tool. Some tools may extract the archive to a folder within your chosen installation location. If this occurs, you can move the contents of the subfolder into the chosen installation location.

2.3.5.2 Creating an Option File

If you need to specify startup options when you run the server, you can indicate them on the command line or place them in an option file. For options that are used every time the server starts, you may find it most convenient to use an option file to specify your MySQL configuration. This is particularly true under the following circumstances:

- The installation or data directory locations are different from the default locations (`C:\Program Files\MySQL\MySQL Server 8.0` and `C:\Program Files\MySQL\MySQL Server 8.0\data`).
- You need to tune the server settings, such as memory, cache, or InnoDB configuration information.

When the MySQL server starts on Windows, it looks for option files in several locations, such as the Windows directory, `C:\`, and the MySQL installation directory (for the full list of locations, see Section 4.2.6, “Using Option Files”). The Windows directory typically is named something like `C:\WINDOWS`. You can determine its exact location from the value of the `WINDIR` environment variable using the following command:

```
C:\> echo %WINDIR%
```

MySQL looks for options in each location first in the `my.ini` file, and then in the `my.cnf` file. However, to avoid confusion, it is best if you use only one file. If your PC uses a boot loader where `C:` is not the boot drive, your only option is to use the `my.ini` file. Whichever option file you use, it must be a plain text file.

Note

When using the MySQL Installer to install MySQL Server, it will create the `my.ini` at the default location, and the user executing MySQL Installer is granted full permissions to this new `my.ini` file.

In other words, be sure that the MySQL Server user has permission to read the `my.ini` file.

You can also make use of the example option files included with your MySQL distribution; see Section 5.1.2, “Server Configuration Defaults”.

An option file can be created and modified with any text editor, such as Notepad. For example, if MySQL is installed in `E:\mysql` and the data directory is in `E:\mydata\data`, you can create an option file containing a `[mysqld]` section to specify values for the `basedir` and `datadir` options:

```
[mysqld]
```

```
# set basedir to your installation path
basedir=E:/mysql
# set datadir to the location of your data directory
datadir=E:/mydata/data
```

Microsoft Windows path names are specified in option files using (forward) slashes rather than backslashes. If you do use backslashes, double them:

```
[mysqld]
# set basedir to your installation path
basedir=E:\\mysql
# set datadir to the location of your data directory
datadir=E:\\mydata\\data
```

The rules for use of backslash in option file values are given in Section 4.2.6, “Using Option Files”.

The Zip Archive does not include a `data` directory. To initialize a MySQL installation by creating the data directory and populating the tables in the mysql system database, initialize MySQL using either `--initialize` or `--initialize-insecure`. For additional information, see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

If you would like to use a data directory in a different location, you should copy the entire contents of the `data` directory to the new location. For example, if you want to use `E:\mydata` as the data directory instead, you must do two things:

1. Move the entire `data` directory and all of its contents from the default location (for example `C:\Program Files\MySQL\MySQL Server 8.0\data`) to `E:\mydata`.
2. Use a `--datadir` option to specify the new data directory location each time you start the server.

2.3.5.3 Selecting a MySQL Server Type

The following table shows the available servers for Windows in MySQL 8.0.

Binary	Description
`mysqld`	Optimized binary with named-pipe support
`mysqld-debug`	Like `mysqld`, but compiled with full debugging and automatic memory allocation checking

All of the preceding binaries are optimized for modern Intel processors, but should work on any Intel i386-class or higher processor.

Each of the servers in a distribution support the same set of storage engines. The `SHOW ENGINES` statement displays which engines a given server supports.

All Windows MySQL 8.0 servers have support for symbolic linking of database directories.

MySQL supports TCP/IP on all Windows platforms. MySQL servers on Windows also support named pipes, if you start the server with the `--enable-named-pipe` option. It is necessary to use this option explicitly because some users have experienced problems with shutting down the MySQL server when named pipes were used. The default is to use TCP/IP regardless of platform because named pipes are slower than TCP/IP in many Windows configurations.

2.3.5.4 Initializing the Data Directory

If you installed MySQL using the `Noinstall` package, no data directory is included. To initialize the data directory, use the instructions at Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

2.3.5.5 Starting the Server for the First Time

This section gives a general overview of starting the MySQL server. The following sections provide more specific information for starting the MySQL server from the command line or as a Windows service.

The information here applies primarily if you installed MySQL using the `Noinstall` version, or if you wish to configure and test MySQL manually rather than with the GUI tools.

Note

The MySQL server will automatically start after using MySQL Installer, and MySQL Notifier can be used to start/stop/restart at any time.

The examples in these sections assume that MySQL is installed under the default location of `C:\Program Files\MySQL\MySQL Server 8.0`. Adjust the path names shown in the examples if you have MySQL installed in a different location.

Clients have two options. They can use TCP/IP, or they can use a named pipe if the server supports named-pipe connections.

MySQL for Windows also supports shared-memory connections if the server is started with the `--shared-memory` option. Clients can connect through shared memory by using the `--protocol=MEMORY` option.

For information about which server binary to run, see Section 2.3.5.3, “Selecting a MySQL Server Type”.

Testing is best done from a command prompt in a console window (or “DOS window”). In this way you can have the server display status messages in the window where they are easy to see. If something is wrong with your configuration, these messages make it easier for you to identify and fix any problems.

Note

The database must be initialized before MySQL can be started. For additional information about the initialization process, see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

To start the server, enter this command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld" --console
```

For a server that includes `InnoDB` support, you should see the messages similar to those following as it starts (the path names and sizes may differ):

```
InnoDB: The first specified datafile c:\ibdata\ibdata1 did not exist:
InnoDB: a new database to be created!
InnoDB: Setting file c:\ibdata\ibdata1 size to 209715200
InnoDB: Database physically writes the file full: wait...
InnoDB: Log file c:\iblogs\ib_logfile0 did not exist: new to be created
InnoDB: Setting log file c:\iblogs\ib_logfile0 size to 31457280
InnoDB: Log file c:\iblogs\ib_logfile1 did not exist: new to be created
InnoDB: Setting log file c:\iblogs\ib_logfile1 size to 31457280
InnoDB: Log file c:\iblogs\ib_logfile2 did not exist: new to be created
InnoDB: Setting log file c:\iblogs\ib_logfile2 size to 31457280
InnoDB: Doublewrite buffer not found: creating new
InnoDB: Doublewrite buffer created
InnoDB: creating foreign key constraint system tables
InnoDB: foreign key constraint system tables created
011024 10:58:25  InnoDB: Started
```

When the server finishes its startup sequence, you should see something like this, which indicates that the server is ready to service client connections:

```
mysqld: ready for connections
Version: '8.0.3'  socket: ''  port: 3306
```

The server continues to write to the console any further diagnostic output it produces. You can open a new console window in which to run client programs.

If you omit the `--console` option, the server writes diagnostic output to the error log in the data directory (`C:\Program Files\MySQL\MySQL Server 8.0\data` by default). The error log is the file with the `.err` extension, and may be set using the `--log-error` option.

Note

The initial `root` account in the MySQL grant tables has no password. After starting the server, you should set up a password for it using the instructions in Section 2.9.4, “Securing the Initial MySQL Account”.

2.3.5.6 Starting MySQL from the Windows Command Line

The MySQL server can be started manually from the command line. This can be done on any version of Windows.

Note

MySQL Notifier can also be used to start/stop/restart the MySQL server.

To start the `mysqld` server from the command line, you should start a console window (or “DOS window”) and enter this command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld"
```

The path to `mysqld` may vary depending on the install location of MySQL on your system.

You can stop the MySQL server by executing this command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqladmin" -u root shutdown
```

Note

If the MySQL `root` user account has a password, you need to invoke `mysqladmin` with the `-p` option and supply the password when prompted.

This command invokes the MySQL administrative utility `mysqladmin` to connect to the server and tell it to shut down. The command connects as the MySQL `root` user, which is the default administrative account in the MySQL grant system.

Note

Users in the MySQL grant system are wholly independent from any login users under Microsoft Windows.

If `mysqld` doesn't start, check the error log to see whether the server wrote any messages there to indicate the cause of the problem. By default, the error log is located in the `C:\Program Files\MySQL\MySQL Server 8.0\data` directory. It is the file with a suffix of `.err`, or may be specified by passing in the `--log-error` option. Alternatively, you can try to start the server with the `--console` option; in this case, the server may display some useful information on the screen that will help solve the problem.

The last option is to start `mysqld` with the `--standalone` and `--debug` options. In this case, `mysqld` writes a log file `C:\mysqld.trace` that should contain the reason why `mysqld` doesn't start. See Section 28.5.3, “The DBUG Package”.

Use `mysqld --verbose --help` to display all the options that `mysqld` supports.

2.3.5.7 Customizing the PATH for MySQL Tools

Warning

You must exercise great care when editing your system `PATH` by hand; accidental deletion or modification of any portion of the existing `PATH` value can leave you with a malfunctioning or even unusable system.

To make it easier to invoke MySQL programs, you can add the path name of the MySQL `bin` directory to your Windows system `PATH` environment variable:

- On the Windows desktop, right-click the **My Computer** icon, and select **Properties**.
- Next select the **Advanced** tab from the **System Properties** menu that appears, and click the **Environment Variables** button.
- Under **System Variables**, select **Path**, and then click the **Edit** button. The **Edit System Variable** dialogue should appear.
- Place your cursor at the end of the text appearing in the space marked **Variable Value**. (Use the **End** key to ensure that your cursor is positioned at the very end of the text in this space.) Then enter the complete path name of your MySQL `bin` directory (for example, `C:\Program Files\MySQL\MySQL Server 8.0\bin`)

  Note

  There must be a semicolon separating this path from any values present in this field.

  Dismiss this dialogue, and each dialogue in turn, by clicking **OK** until all of the dialogues that were opened have been dismissed. The new `PATH` value should now be available to any new command shell you open, allowing you to invoke any MySQL executable program by typing its name at the DOS prompt from any directory on the system, without having to supply the path. This includes the servers, the `mysql` client, and all MySQL command-line utilities such as `mysqladmin` and `mysqldump`.

You should not add the MySQL `bin` directory to your Windows `PATH` if you are running multiple MySQL servers on the same machine.

2.3.5.8 Starting MySQL as a Windows Service

On Windows, the recommended way to run MySQL is to install it as a Windows service, so that MySQL starts and stops automatically when Windows starts and stops. A MySQL server installed as a service can also be controlled from the command line using `NET` commands, or with the graphical `Services` utility. Generally, to install MySQL as a Windows service you should be logged in using an account that has administrator rights.

Note

MySQL Notifier can also be used to monitor the status of the MySQL service.

The `Services` utility (the Windows `Service Control Manager`) can be found in the Windows Control Panel (under **Administrative Tools** on Windows Vista, and Server 2003). To avoid conflicts, it is advisable to close the `Services` utility while performing server installation or removal operations from the command line.

Installing the service

Before installing MySQL as a Windows service, you should first stop the current server if it is running by using the following command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqladmin"
        -u root shutdown
```

Note

If the MySQL `root` user account has a password, you need to invoke `mysqladmin` with the `-p` option and supply the password when prompted.

This command invokes the MySQL administrative utility `mysqladmin` to connect to the server and tell it to shut down. The command connects as the MySQL `root` user, which is the default administrative account in the MySQL grant system.

Note

Users in the MySQL grant system are wholly independent from any login users under Windows.

Install the server as a service using this command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld" --install
```

The service-installation command does not start the server. Instructions for that are given later in this section.

To make it easier to invoke MySQL programs, you can add the path name of the MySQL `bin` directory to your Windows system `PATH` environment variable:

- On the Windows desktop, right-click the **My Computer** icon, and select **Properties**.
- Next select the **Advanced** tab from the **System Properties** menu that appears, and click the **Environment Variables** button.
- Under **System Variables**, select **Path**, and then click the **Edit** button. The **Edit System Variable** dialogue should appear.
- Place your cursor at the end of the text appearing in the space marked **Variable Value**. (Use the **End** key to ensure that your cursor is positioned at the very end of the text in this space.) Then enter the complete path name of your MySQL `bin` directory (for example, `C:\Program Files\MySQL\MySQL Server 8.0\bin`), and there should be a semicolon separating this path from any values present in this field. Dismiss this dialogue, and each dialogue in turn, by clicking **OK** until all of the dialogues that were opened have been dismissed. You should now be able to invoke any MySQL executable program by typing its name at the DOS prompt from any directory on the system, without having to supply the path. This includes the servers, the `mysql` client, and all MySQL command-line utilities such as `mysqladmin` and `mysqldump`.

  You should not add the MySQL `bin` directory to your Windows `PATH` if you are running multiple MySQL servers on the same machine.

Warning

You must exercise great care when editing your system `PATH` by hand; accidental deletion or modification of any portion of the existing `PATH` value can leave you with a malfunctioning or even unusable system.

The following additional arguments can be used when installing the service:

- You can specify a service name immediately following the `--install` option. The default service name is `MySQL`.
- If a service name is given, it can be followed by a single option. By convention, this should be `--defaults-file=file_name` to specify the name of an option file from which the server should read options when it starts.

  The use of a single option other than `--defaults-file` is possible but discouraged. `--defaults-file` is more flexible because it enables you to specify multiple startup options for the server by placing them in the named option file.
- You can also specify a `--local-service` option following the service name. This causes the server to run using the `LocalService` Windows account that has limited system privileges. If both `--defaults-file` and `--local-service` are given following the service name, they can be in any order.

For a MySQL server that is installed as a Windows service, the following rules determine the service name and option files that the server uses:

- If the service-installation command specifies no service name or the default service name (`MySQL`) following the `--install` option, the server uses the service name of `MySQL` and reads options from the `[mysqld]` group in the standard option files.
- If the service-installation command specifies a service name other than `MySQL` following the `--install` option, the server uses that service name. It reads options from the `[mysqld]` group and the group that has the same name as the service in the standard option files. This enables you to use the `[mysqld]` group for options that should be used by all MySQL services, and an option group with the service name for use by the server installed with that service name.
- If the service-installation command specifies a `--defaults-file` option after the service name, the server reads options the same way as described in the previous item, except that it reads options only from the named file and ignores the standard option files.

As a more complex example, consider the following command:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld"
          --install MySQL --defaults-file=C:\my-opts.cnf
```

Here, the default service name (`MySQL`) is given after the `--install` option. If no `--defaults-file` option had been given, this command would have the effect of causing the server to read the `[mysqld]` group from the standard option files. However, because the `--defaults-file` option is present, the server reads options from the `[mysqld]` option group, and only from the named file.

Note

On Windows, if the server is started with the `--defaults-file` and `--install` options, `--install` must be first. Otherwise, `mysqld.exe` will attempt to start the MySQL server.

You can also specify options as Start parameters in the Windows `Services` utility before you start the MySQL service.

Finally, before trying to start the MySQL service, make sure the user variables `%TEMP%` and `%TMP%` (and also `%TMPDIR%`, if it has ever been set) for the system user who is to run the service are pointing to a folder to which the user has write access. The default user for running the MySQL service is `LocalSystem`, and the default value for its `%TEMP%` and `%TMP%` is `C:\Windows\Temp`, a directory `LocalSystem` has write access to by default. However, if there are any changes to that default setup (for example, changes to the user who runs the service or to the mentioned user variables, or the `--tmpdir` option has been used to put the temporary directory somewhere else), the MySQL service might fail to run because write access to the temporary directory has not been granted to the proper user.

Starting the service

Once a MySQL server has been installed as a service, Windows starts the service automatically whenever Windows starts. The service also can be started immediately from the `Services` utility, or by using a `NET START MySQL` command. The `NET` command is not case sensitive.

When run as a service, `mysqld` has no access to a console window, so no messages can be seen there. If `mysqld` does not start, check the error log to see whether the server wrote any messages there to indicate the cause of the problem. The error log is located in the MySQL data directory (for example, `C:\Program Files\MySQL\MySQL Server 8.0\data`). It is the file with a suffix of `.err`.

When a MySQL server has been installed as a service, and the service is running, Windows stops the service automatically when Windows shuts down. The server also can be stopped manually by using the `Services` utility, the `NET STOP MySQL` command, or the `mysqladmin shutdown` command.

You also have the choice of installing the server as a manual service if you do not wish for the service to be started automatically during the boot process. To do this, use the `--install-manual` option rather than the `--install` option:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld" --install-manual
```

Removing the service

To remove a server that is installed as a service, first stop it if it is running by executing `NET STOP MySQL`. Then use the `--remove` option to remove it:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld" --remove
```

If `mysqld` is not running as a service, you can start it from the command line. For instructions, see Section 2.3.5.6, “Starting MySQL from the Windows Command Line”.

If you encounter difficulties during installation, see Section 2.3.6, “Troubleshooting a Microsoft Windows MySQL Server Installation”.

For more information about stopping or removing a MySQL Windows service, see Section 5.7.2.2, “Starting Multiple MySQL Instances as Windows Services”.

2.3.5.9 Testing The MySQL Installation

You can test whether the MySQL server is working by executing any of the following commands:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqlshow"
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqlshow" -u root mysql
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqladmin" version status proc
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysql" test
```

If `mysqld` is slow to respond to TCP/IP connections from client programs, there is probably a problem with your DNS. In this case, start `mysqld` with the `--skip-name-resolve` option and use only `localhost` and IP addresses in the `Host` column of the MySQL grant tables. (Be sure that an account exists that specifies an IP address or you may not be able to connect.)

You can force a MySQL client to use a named-pipe connection rather than TCP/IP by specifying the `--pipe` or `--protocol=PIPE` option, or by specifying `.` (period) as the host name. Use the `--socket` option to specify the name of the pipe if you do not want to use the default pipe name.

If you have set a password for the `root` account, deleted the anonymous account, or created a new user account, then to connect to the MySQL server you must use the appropriate `-u` and `-p` options with the commands shown previously. See Section 4.2.2, “Connecting to the MySQL Server”.

For more information about `mysqlshow`, see Section 4.5.7, “`mysqlshow` — Display Database, Table, and Column Information”.

2.3.6 Troubleshooting a Microsoft Windows MySQL Server Installation

When installing and running MySQL for the first time, you may encounter certain errors that prevent the MySQL server from starting. This section helps you diagnose and correct some of these errors.

Your first resource when troubleshooting server issues is the error log. The MySQL server uses the error log to record information relevant to the error that prevents the server from starting. The error log is located in the data directory specified in your `my.ini` file. The default data directory location is `C:\Program Files\MySQL\MySQL Server 8.0\data`, or `C:\ProgramData\Mysql` on Windows 7 and Windows Server 2008. The `C:\ProgramData` directory is hidden by default. You need to change your folder options to see the directory and contents. For more information on the error log and understanding the content, see Section 5.4.2, “The Error Log”.

For information regarding possible errors, also consult the console messages displayed when the MySQL service is starting. Use the `NET START MySQL` command from the command line after installing `mysqld` as a service to see any error messages regarding the starting of the MySQL server as a service. See Section 2.3.5.8, “Starting MySQL as a Windows Service”.

The following examples show other common error messages you might encounter when installing MySQL and starting the server for the first time:

- If the MySQL server cannot find the `mysql` privileges database or other critical files, it displays these messages:

```
System error 1067 has occurred.
Fatal error: Can't open and lock privilege tables:
Table 'mysql.user' doesn't exist
```

  These messages often occur when the MySQL base or data directories are installed in different locations than the default locations (`C:\Program Files\MySQL\MySQL Server 8.0` and `C:\Program Files\MySQL\MySQL Server 8.0\data`, respectively).

  This situation can occur when MySQL is upgraded and installed to a new location, but the configuration file is not updated to reflect the new location. In addition, old and new configuration files might conflict. Be sure to delete or rename any old configuration files when upgrading MySQL.

  If you have installed MySQL to a directory other than `C:\Program Files\MySQL\MySQL Server 8.0`, ensure that the MySQL server is aware of this through the use of a configuration (`my.ini`) file. Put the `my.ini` file in your Windows directory, typically `C:\WINDOWS`. To determine its exact location from the value of the `WINDIR` environment variable, issue the following command from the command prompt:

```
C:\> echo %WINDIR%
```

  You can create or modify an option file with any text editor, such as Notepad. For example, if MySQL is installed in `E:\mysql` and the data directory is `D:\MySQLdata`, you can create the option file and set up a `[mysqld]` section to specify values for the `basedir` and `datadir` options:

```
[mysqld]
# set basedir to your installation path
basedir=E:/mysql
# set datadir to the location of your data directory
datadir=D:/MySQLdata
```

  Microsoft Windows path names are specified in option files using (forward) slashes rather than backslashes. If you do use backslashes, double them:

```
[mysqld]
# set basedir to your installation path
```

```
basedir=C:\\Program Files\\MySQL\\MySQL Server 8.0
# set datadir to the location of your data directory
datadir=D:\\MySQLdata
```

The rules for use of backslash in option file values are given in Section 4.2.6, “Using Option Files”.

If you change the `datadir` value in your MySQL configuration file, you must move the contents of the existing MySQL data directory before restarting the MySQL server.

See Section 2.3.5.2, “Creating an Option File”.

- If you reinstall or upgrade MySQL without first stopping and removing the existing MySQL service and install MySQL using the MySQL Installer, you might see this error:

```
Error: Cannot create Windows service for MySql. Error: 0
```

This occurs when the Configuration Wizard tries to install the service and finds an existing service with the same name.

One solution to this problem is to choose a service name other than `mysql` when using the configuration wizard. This enables the new service to be installed correctly, but leaves the outdated service in place. Although this is harmless, it is best to remove old services that are no longer in use.

To permanently remove the old `mysql` service, execute the following command as a user with administrative privileges, on the command line:

```
C:\> sc delete mysql
[SC] DeleteService SUCCESS
```

If the `sc` utility is not available for your version of Windows, download the `delsrv` utility from http://www.microsoft.com/windows2000/techinfo/reskit/tools/existing/delsrv-o.asp and use the `delsrv mysql` syntax.

2.3.7 Windows Postinstallation Procedures

GUI tools exist that perform most of the tasks described in this section, including:

- MySQL Installer: Used to install and upgrade MySQL products.
- MySQL Workbench: Manages the MySQL server and edits SQL statements.
- MySQL Notifier: Starts, stops, or restarts the MySQL server, and monitors its status.
- MySQL for Excel: Edits MySQL data with Microsoft Excel.

If necessary, initialize the data directory and create the MySQL grant tables. Windows installation operations performed by MySQL Installer initialize the data directory automatically. For installation from a Zip package, you can initialize the data directory as described at Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

Regarding passwords, if you installed MySQL using the MySQL Installer, you may have already assigned a passwords to the initial `root` account. (See Section 2.3.3, “MySQL Installer for Windows”.) Otherwise, use the password-assignment procedure given in Section 2.9.4, “Securing the Initial MySQL Account”.

Before assigning passwords, you might want to try running some client programs to make sure that you can connect to the server and that it is operating properly. Make sure that the server is running (see Section 2.3.5.5, “Starting the Server for the First Time”). You can also set up a MySQL service that runs automatically when Windows starts (see Section 2.3.5.8, “Starting MySQL as a Windows Service”).

These instructions assume that your current location is the MySQL installation directory and that it has a `bin` subdirectory containing the MySQL programs used here. If that is not true, adjust the command path names accordingly.

If you installed MySQL using MySQL Installer (see Section 2.3.3, “MySQL Installer for Windows”), the default installation directory is `C:\Program Files\MySQL\MySQL Server 8.0`:

```
C:\> cd "C:\Program Files\MySQL\MySQL Server 8.0"
```

A common installation location for installation from a Zip package is `C:\mysql`:

```
C:\> cd C:\mysql
```

Alternatively, add the `bin` directory to your `PATH` environment variable setting. That enables your command interpreter to find MySQL programs properly, so that you can run a program by typing only its name, not its path name. See Section 2.3.5.7, “Customizing the PATH for MySQL Tools”.

With the server running, issue the following commands to verify that you can retrieve information from the server. The output should be similar to that shown here.

Use `mysqlshow` to see what databases exist:

```
C:\> bin\mysqlshow
+--------------------+
|     Databases      |
+--------------------+
| information_schema |
| mysql              |
| performance_schema |
| sys                |
+--------------------+
```

The list of installed databases may vary, but will always include the minimum of `mysql` and `information_schema`.

The preceding command (and commands for other MySQL programs such as `mysql`) may not work if the correct MySQL account does not exist. For example, the program may fail with an error, or you may not be able to view all databases. If you installed MySQL using MySQL Installer, the `root` user will have been created automatically with the password you supplied. In this case, you should use the `-u root` and `-p` options. (You must use those options if you have already secured the initial MySQL accounts.) With `-p`, the client program prompts for the `root` password. For example:

```
C:\> bin\mysqlshow -u root -p
Enter password: (enter root password here)
+--------------------+
|     Databases      |
+--------------------+
| information_schema |
| mysql              |
| performance_schema |
| sys                |
+--------------------+
```

If you specify a database name, `mysqlshow` displays a list of the tables within the database:

```
C:\> bin\mysqlshow mysql
Database: mysql
+---------------------------+
|          Tables           |
+---------------------------+
| columns_priv              |
| db                        |
```

```
| engine_cost               |
| event                     |
| func                      |
| general_log               |
| gtid_executed             |
| help_category             |
| help_keyword              |
| help_relation             |
| help_topic                |
| innodb_index_stats        |
| innodb_table_stats        |
| ndb_binlog_index          |
| plugin                    |
| proc                      |
| procs_priv                |
| proxies_priv              |
| server_cost               |
| servers                   |
| slave_master_info         |
| slave_relay_log_info      |
| slave_worker_info         |
| slow_log                  |
| tables_priv               |
| time_zone                 |
| time_zone_leap_second     |
| time_zone_name            |
| time_zone_transition      |
| time_zone_transition_type |
| user                      |
+---------------------------+
```

Use the `mysql` program to select information from a table in the `mysql` database:

```
C:\> bin\mysql -e "SELECT User, Host, plugin FROM mysql.user" mysql
+------+-----------+-----------------------+
| User | Host      | plugin                |
+------+-----------+-----------------------+
| root | localhost | mysql_native_password |
+------+-----------+-----------------------+
```

For more information about `mysql` and `mysqlshow`, see Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”, and Section 4.5.7, “`mysqlshow` — Display Database, Table, and Column Information”.

2.3.8 Upgrading MySQL on Windows

To upgrade MySQL on Windows, follow these steps:

1. Review Section 2.10.1, “Upgrading MySQL”, for additional information on upgrading MySQL that is not specific to Windows.
2. Always back up your current MySQL installation before performing an upgrade. See Section 7.2, “Database Backup Methods”.
3. Download the latest Windows distribution of MySQL from http://dev.mysql.com/downloads/.
4. Before upgrading MySQL, stop the server. If the server is installed as a service, stop the service with the following command from the command prompt:

   ```
   C:\> NET STOP MySQL
   ```

   If you are not running the MySQL server as a service, use `mysqladmin` to stop it. For example, before upgrading from MySQL 5.7 to 8.0, use `mysqladmin` from MySQL 5.7 as follows:

   ```
   C:\> "C:\Program Files\MySQL\MySQL Server 5.7\bin\mysqladmin" -u root shutdown
   ```

Note

If the MySQL `root` user account has a password, invoke `mysqladmin` with the `-p` option and enter the password when prompted.

5. Before upgrading to MySQL 8.0 from a version previous to 4.1.5, or from a version of MySQL installed from a Zip archive to a version of MySQL installed with the MySQL Installation Wizard, you must first manually remove the previous installation and MySQL service (if the server is installed as a service).

   To remove the MySQL service, use the following command:

   ```
   C:\> C:\mysql\bin\mysqld --remove
   ```

   If you do not remove the existing service, the MySQL Installation Wizard may fail to properly install the new MySQL service.

6. If you are using the MySQL Installer, start it as described in Section 2.3.3, “MySQL Installer for Windows”.

7. If you are upgrading MySQL from a Zip archive, extract the archive. You may either overwrite your existing MySQL installation (usually located at `C:\mysql`), or install it into a different directory, such as `C:\mysql5`. Overwriting the existing installation is recommended. However, for upgrades (as opposed to installing for the first time), you must remove the data directory from your existing MySQL installation to avoid replacing your current data files. To do so, follow these steps:

   a. Unzip the Zip archive in some location other than your current MySQL installation

   b. Remove the data directory

   c. Rezip the Zip archive

   d. Unzip the modified Zip archive on top of your existing installation

   Alternatively:

   a. Unzip the Zip archive in some location other than your current MySQL installation

   b. Remove the data directory

   c. Move the data directory from the current MySQL installation to the location of the just-removed data directory

   d. Remove the current MySQL installation

   e. Move the unzipped installation to the location of the just-removed installation

8. If you were running MySQL as a Windows service and you had to remove the service earlier in this procedure, reinstall the service. (See Section 2.3.5.8, “Starting MySQL as a Windows Service”.)

9. Restart the server. For example, use `NET START MySQL` if you run MySQL as a service, or invoke `mysqld` directly otherwise.

10. As Administrator, run `mysql_upgrade` to check your tables, attempt to repair them if necessary, and update your grant tables if they have changed so that you can take advantage of any new capabilities. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

11. If you encounter errors, see Section 2.3.6, “Troubleshooting a Microsoft Windows MySQL Server Installation”.

2.4 Installing MySQL on OS X

For a list of OS X versions that the MySQL server supports, see http://www.mysql.com/support/supportedplatforms/database.html.

MySQL for OS X is available in a number of different forms:

- Native Package Installer, which uses the native OS X installer (DMG) to walk you through the installation of MySQL. For more information, see Section 2.4.2, “Installing MySQL on OS X Using Native Packages”. You can use the package installer with OS X. The user you use to perform the installation must have administrator privileges.
- Compressed TAR archive, which uses a file packaged using the Unix `tar` and `gzip` commands. To use this method, you will need to open a `Terminal` window. You do not need administrator privileges using this method, as you can install the MySQL server anywhere using this method. For more information on using this method, you can use the generic instructions for using a tarball, Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

  In addition to the core installation, the Package Installer also includes Section 2.4.3, “Installing a MySQL Launch Daemon” and Section 2.4.4, “Installing and Using the MySQL Preference Pane”, both of which simplify the management of your installation.

For additional information on using MySQL on OS X, see Section 2.4.1, “General Notes on Installing MySQL on OS X”.

2.4.1 General Notes on Installing MySQL on OS X

You should keep the following issues and notes in mind:

- A launchd daemon is installed, and it includes MySQL configuration options. Consider editing it if needed, see the documentation below for additional information. Also, macOS 10.10 removed startup item support in favor of launchd daemons. The optional MySQL preference pane under macOS **System Preferences** uses the launchd daemon.
- You may need (or want) to create a specific `mysql` user to own the MySQL directory and data. You can do this through the `Directory Utility`, and the `mysql` user should already exist. For use in single user mode, an entry for `_mysql` (note the underscore prefix) should already exist within the system `/etc/passwd` file.
- Because the MySQL package installer installs the MySQL contents into a version and platform specific directory, you can use this to upgrade and migrate your database between versions. You will need to either copy the `data` directory from the old version to the new version, or alternatively specify an alternative `datadir` value to set location of the data directory. By default, the MySQL directories are installed under `/usr/local/`.
- You might want to add aliases to your shell's resource file to make it easier to access commonly used programs such as `mysql` and `mysqladmin` from the command line. The syntax for `bash` is:

  ```
  alias mysql=/usr/local/mysql/bin/mysql
  alias mysqladmin=/usr/local/mysql/bin/mysqladmin
  ```

  For `tcsh`, use:

  ```
  alias mysql /usr/local/mysql/bin/mysql
  alias mysqladmin /usr/local/mysql/bin/mysqladmin
  ```

  Even better, add `/usr/local/mysql/bin` to your `PATH` environment variable. You can do this by modifying the appropriate startup file for your shell. For more information, see Section 4.2.1, “Invoking MySQL Programs”.
- After you have copied over the MySQL database files from the previous installation and have successfully started the new server, you should consider removing the old installation files to save

disk space. Additionally, you should also remove older versions of the Package Receipt directories located in `/Library/Receipts/mysql-VERSION.pkg`.

- Prior to OS X 10.7, MySQL server was bundled with OS X Server.

2.4.2 Installing MySQL on OS X Using Native Packages

The package is located inside a disk image (`.dmg`) file that you first need to mount by double-clicking its icon in the Finder. It should then mount the image and display its contents.

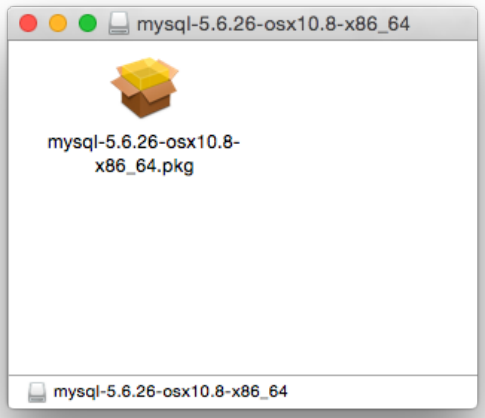
Note

Before proceeding with the installation, be sure to stop all running MySQL server instances by using either the MySQL Manager Application (on OS X Server), the preference pane, or `mysqladmin shutdown` on the command line.

To install MySQL using the package installer:

1. Download the disk image (`.dmg`) file (the community version is available here) that contains the MySQL package installer. Double-click the file to mount the disk image and see its contents.

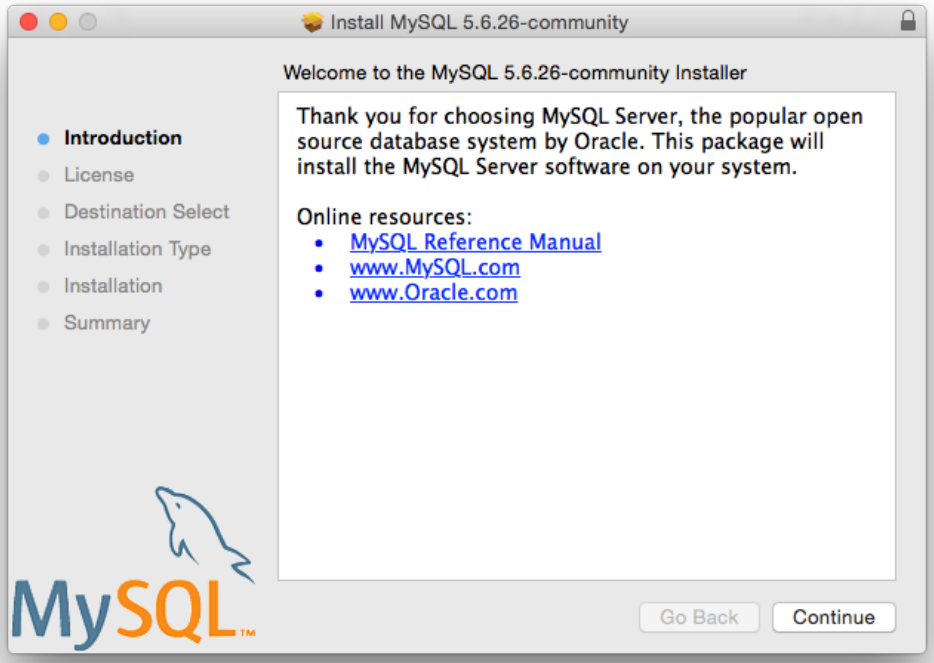
Figure 2.22 MySQL Package Installer: DMG Contents



2. Double-click the MySQL installer package. It will be named according to the version of MySQL you have downloaded. For example, if you have downloaded MySQL server 8.0.3, double-click `mysql-8.0.3-osx-10.11-x86_64.pkg`.

3. You will be presented with the opening installer dialog. Click **Continue** to begin installation.

Figure 2.23 MySQL Package Installer: Introduction



4. If you have downloaded the community version of MySQL, you will be shown a copy of the relevant GNU General Public License. Click **Continue** and then **Agree** to continue.

5. From the **Installation Type** page you can either click **Install** to execute the installation wizard using all defaults, click **Customize** to alter which components to install (MySQL server, Preference Pane, Launchd Support -- all enabled by default), or click **Change Installation Location** to change the type of installation, if available.

Figure 2.24 MySQL Package Installer: Installation Type

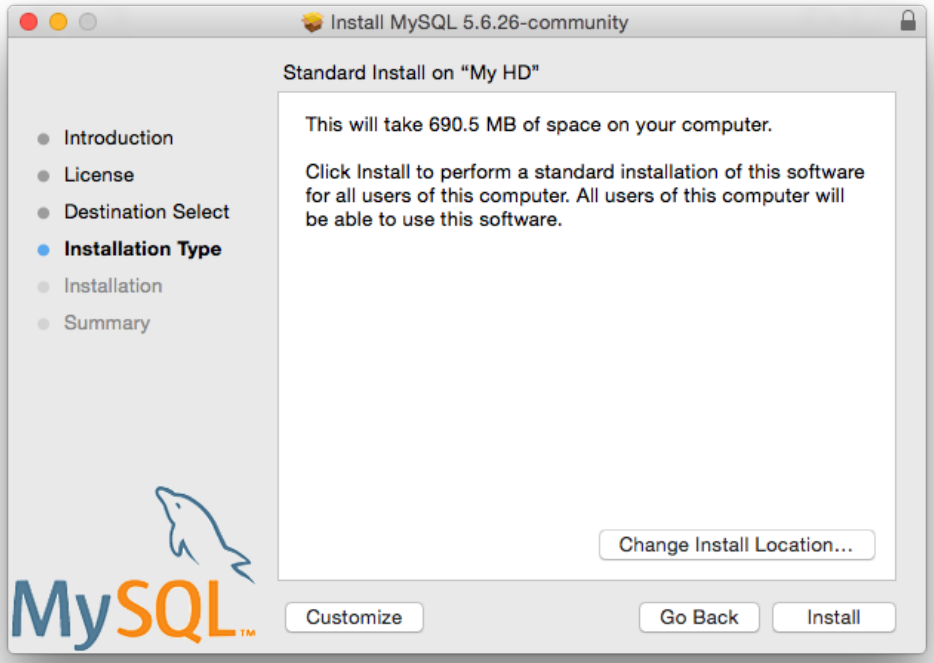


Figure 2.25 MySQL Package Installer: Destination Select (Change Installation Location)

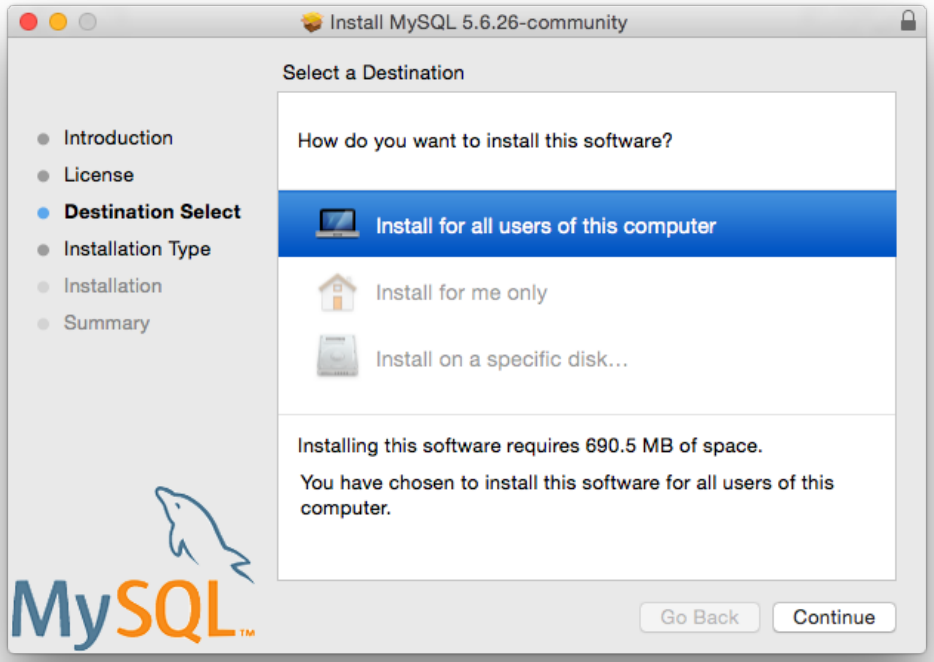
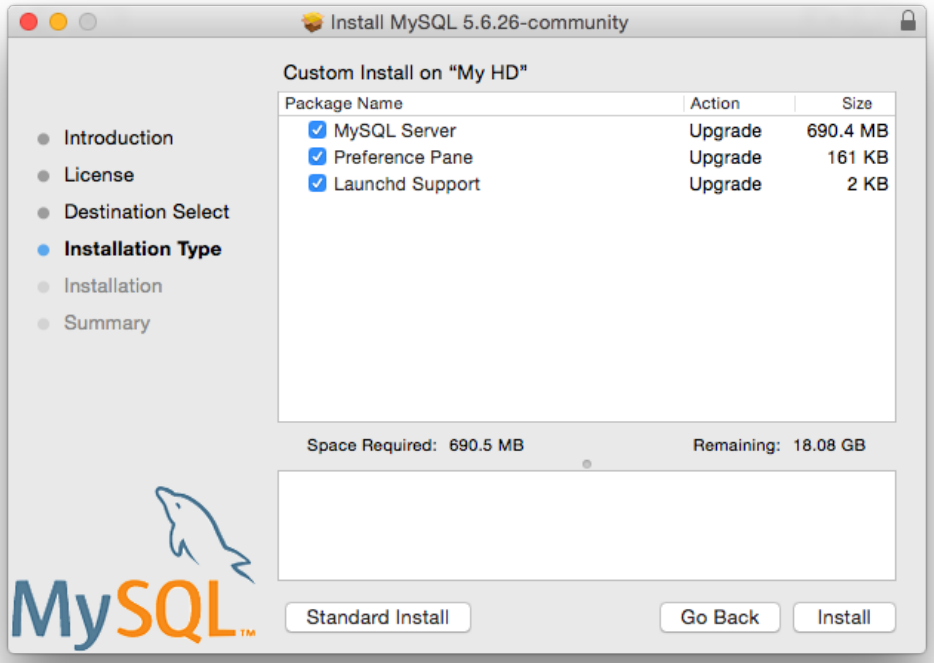
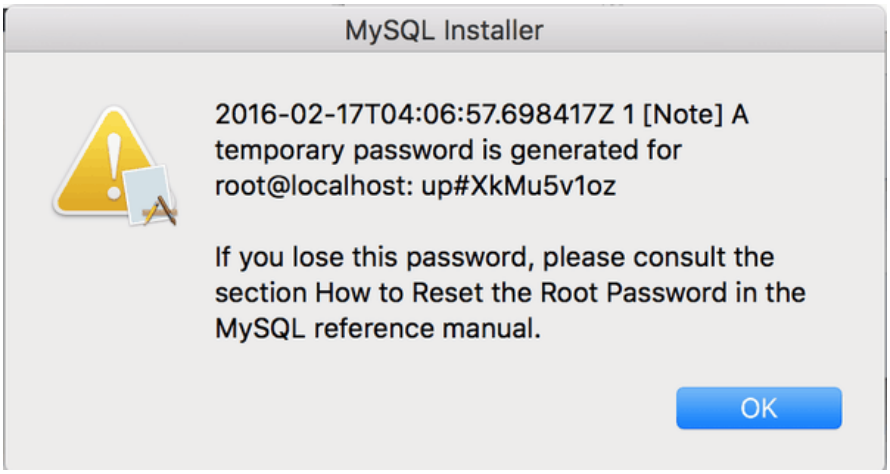


Figure 2.26 MySQL Package Installer: Customize



6. Click **Install** to begin the installation process.
7. Once the installation has been completed successfully, you will be provided with your temporary root password. This cannot be recovered, so you must save this password. For example:

Figure 2.27 MySQL Package Installer: Temporary Root Password

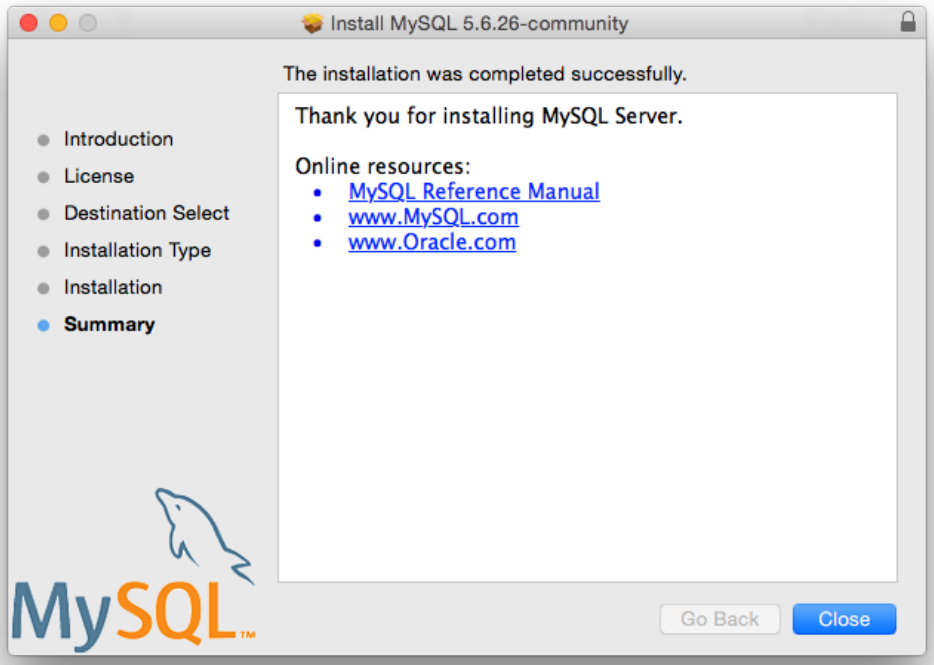


Note

After logging into MySQL using this temporary password, MySQL will expire this password and require you to create a new password.

8. Next is an **Install Succeeded** message with a short summary. **Close** the wizard.

Figure 2.28 MySQL Package Installer: Summary



MySQL server is now installed, but it is not loaded (or started) by default. Use either launchctl from the command line, or start MySQL by clicking "Start" using the MySQL preference pane. For additional information, see Section 2.4.3, “Installing a MySQL Launch Daemon”, and Section 2.4.4, “Installing and Using the MySQL Preference Pane”. Use the MySQL Preference Pane or launchd to configure MySQL to automatically start at bootup.

When installing using the package installer, the files are installed into a directory within `/usr/local` matching the name of the installation version and platform. For example, the installer file `mysql-8.0.3-osx10.11-x86_64.dmg` installs MySQL into `/usr/local/mysql-8.0.3-osx10.11-x86_64/` . The following table shows the layout of the installation directory.

Table 2.6 MySQL Installation Layout on OS X

Directory	Contents of Directory
`bin`	`mysqld` server, client and utility programs
`data`	Log files, databases
`docs`	Helper documents, like the Release Notes and build information
`include`	Include (header) files
`lib`	Libraries
`man`	Unix manual pages
`mysql-test`	MySQL test suite
`share`	Miscellaneous support files, including error messages, sample configuration files, SQL for database installation
`support-files`	Scripts and sample configuration files

Directory	Contents of Directory
`/tmp/mysql.sock`	Location of the MySQL Unix socket

During the package installer process, a symbolic link from `/usr/local/mysql` to the version/platform specific directory created during installation will be created automatically.

2.4.3 Installing a MySQL Launch Daemon

OS X uses launch daemons to automatically start, stop, and manage processes and applications such as MySQL.

By default, the installation package (DMG) on OS X installs a launchd file named `/Library/LaunchDaemons/com.oracle.oss.mysql.mysqld.plist` that contains a plist definition similar to:

```
<?xml version="1.0" encoding="utf-8"?>
<!DOCTYPE plist PUBLIC "-//Apple Computer//DTD PLIST 1.0//EN" "http://www.apple.com/DTDs/PropertyList-1
<plist version="1.0">
<dict>
    <key>Label</key>            <string>com.oracle.oss.mysql.mysqld</string>
    <key>ProcessType</key>      <string>Interactive</string>
    <key>Disabled</key>         <false/>
    <key>RunAtLoad</key>        <true/>
    <key>KeepAlive</key>        <true/>
    <key>SessionCreate</key>    <true/>
    <key>LaunchOnlyOnce</key>   <false/>
    <key>UserName</key>         <string>_mysql</string>
    <key>GroupName</key>        <string>_mysql</string>
    <key>ExitTimeOut</key>      <integer>600</integer>
    <key>Program</key>          <string>/usr/local/mysql/bin/mysqld</string>
    <key>ProgramArguments</key>
        <array>
            <string>/usr/local/mysql/bin/mysqld</string>
            <string>--user=_mysql</string>
            <string>--basedir=/usr/local/mysql</string>
            <string>--datadir=/usr/local/mysql/data</string>
            <string>--plugin-dir=/usr/local/mysql/lib/plugin</string>
            <string>--log-error=/usr/local/mysql/data/mysqld.local.err</string>
            <string>--pid-file=/usr/local/mysql/data/mysqld.local.pid</string>
        </array>
    <key>WorkingDirectory</key>  <string>/usr/local/mysql</string>
</dict>
</plist>
```

Note

Some users report that adding a plist DOCTYPE declaration causes the launchd operation to fail, despite it passing the lint check. We suspect it's a copy-n-paste error. The md5 checksum of a file containing the above snippet is *24710a27dc7a28fb7ee6d825129cd3cf*.

To enable the launchd service, you can either:

- Click **Start MySQL Server** from the MySQL preference pane.

Figure 2.29 MySQL Preference Pane: Location

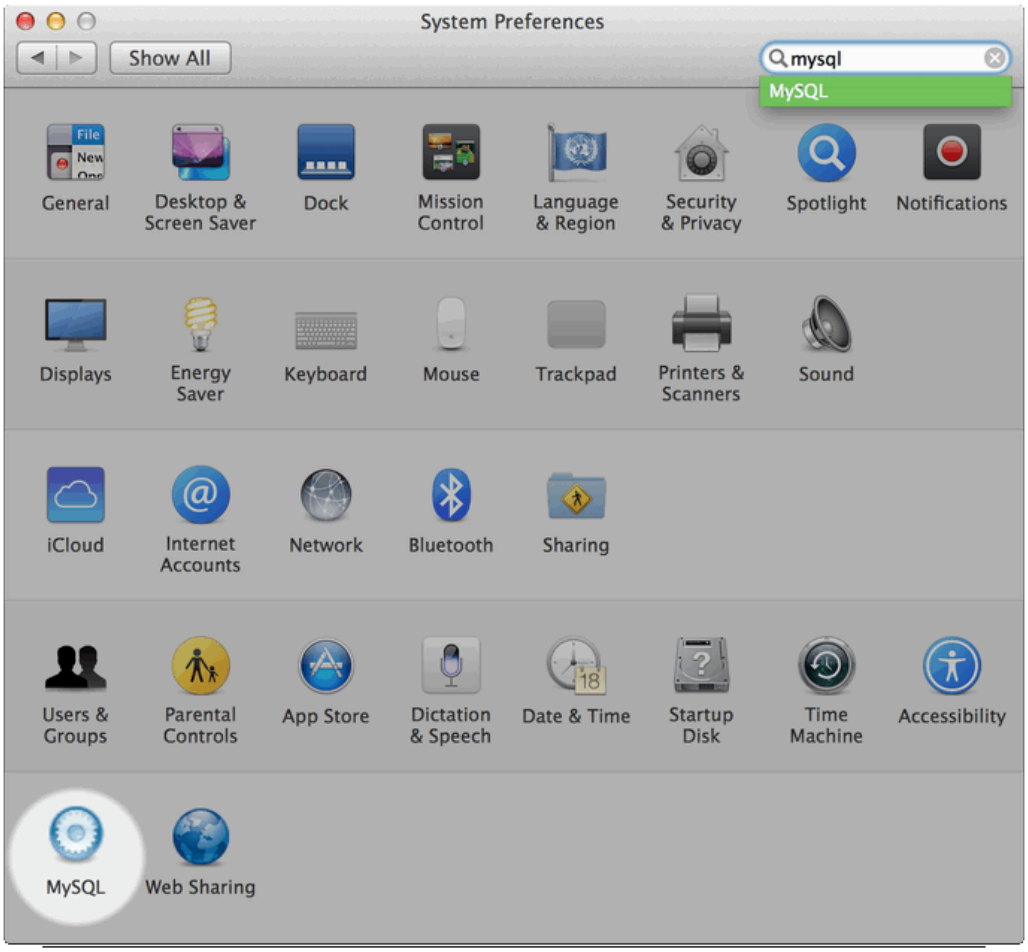
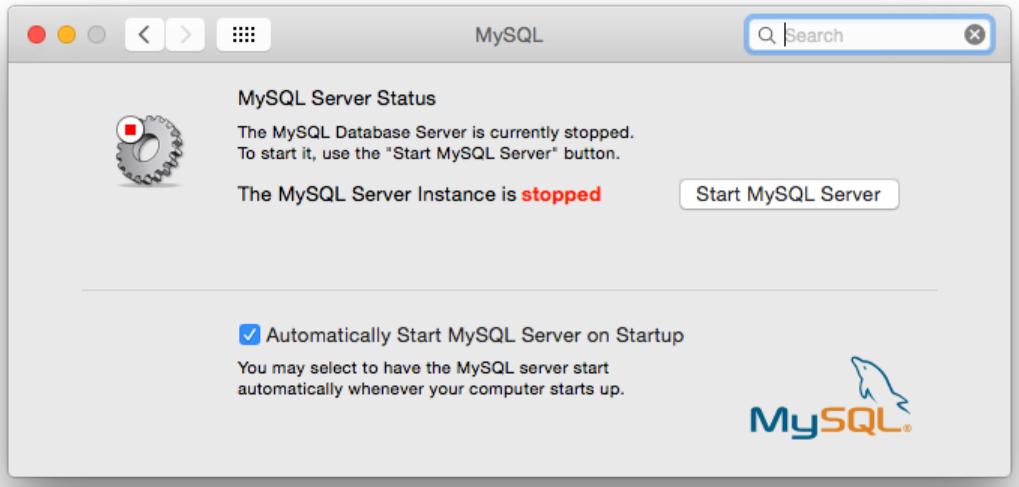


Figure 2.30 MySQL Preference Pane: Usage



- Or, manually load the launchd file.

```
shell> cd /Library/LaunchDaemons
shell> sudo launchctl load -F com.oracle.oss.mysql.mysqld.plist
```

- To configure MySQL to automatically start at bootup, you can:

```
shell> sudo launchctl load -w com.oracle.oss.mysql.mysqld.plist
```

Note

When upgrading MySQL server, the launchd installation process will remove the old startup items that were installed with MySQL server 5.7.7 and below.

Also, upgrading will replace your existing launchd file of the same name.

Additional launchd related information:

- The plist entries override `my.cnf` entries, because they are passed in as command line arguments. For additional information about passing in program options, see Section 4.2.3, “Specifying Program Options”.
- The **ProgramArguments** section defines the command line options that are passed into the program, which is the `mysqld` binary in this case.
- The default plist definition is written with less sophisticated use cases in mind. For more complicated setups, you may want to remove some of the arguments and instead rely on a MySQL configuration file, such as `my.cnf`.
- If you edit the plist file, then uncheck the installer option when reinstalling or upgrading MySQL. Otherwise, your edited plist file will be overwritten, and all edits will be lost.

Because the default plist definition defines several **ProgramArguments**, you might remove most of these arguments and instead rely upon your `my.cnf` MySQL configuration file to define them. For example:

```
<?xml version="1.0" encoding="utf-8"?>
<!DOCTYPE plist PUBLIC "-//Apple Computer//DTD PLIST 1.0//EN" "http://www.apple.com/DTDs/PropertyList-1.0.d
<plist version="1.0">
<dict>
    <key>Label</key>             <string>com.oracle.oss.mysql.mysqld</string>
    <key>ProcessType</key>       <string>Interactive</string>
    <key>Disabled</key>          <false/>
    <key>RunAtLoad</key>         <true/>
    <key>KeepAlive</key>         <true/>
    <key>SessionCreate</key>     <true/>
    <key>LaunchOnlyOnce</key>    <false/>
    <key>UserName</key>          <string>_mysql</string>
    <key>GroupName</key>         <string>_mysql</string>
    <key>ExitTimeOut</key>       <integer>600</integer>
    <key>Program</key>           <string>/usr/local/mysql/bin/mysqld</string>
    <key>WorkingDirectory</key>  <string>/usr/local/mysql</string>
    <key>ProgramArguments</key>
        <array>
            <string>/usr/local/mysql/bin/mysqld</string>
            <string>--user=_mysql</string>
        </array>
</dict>
</plist>
```

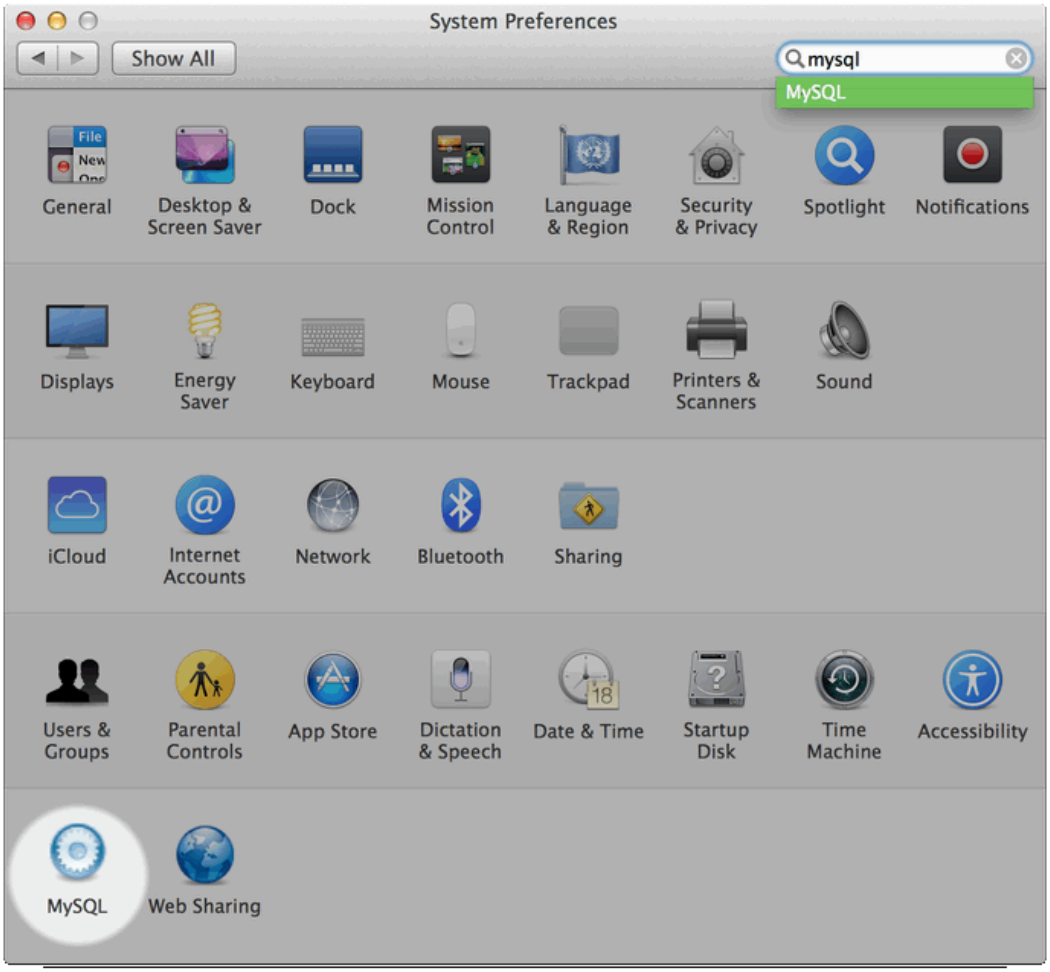
In this case, the `basedir`, `datadir`, `plugin_dir`, `log_error`, and `pid_file` options were removed from the plist definition, and then you you might define them in `my.cnf`.

2.4.4 Installing and Using the MySQL Preference Pane

The MySQL Installation Package includes a MySQL preference pane that enables you to start, stop, and control automated startup during boot of your MySQL installation.

This preference pane is installed by default, and is listed under your system's *System Preferences* window.

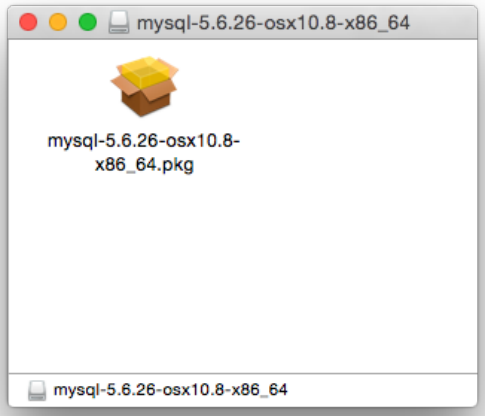
Figure 2.31 MySQL Preference Pane: Location



To install the MySQL Preference Pane:

1. Download the disk image (`.dmg`) file (the community version is available here) that contains the MySQL package installer. Double-click the file to mount the disk image and see its contents.

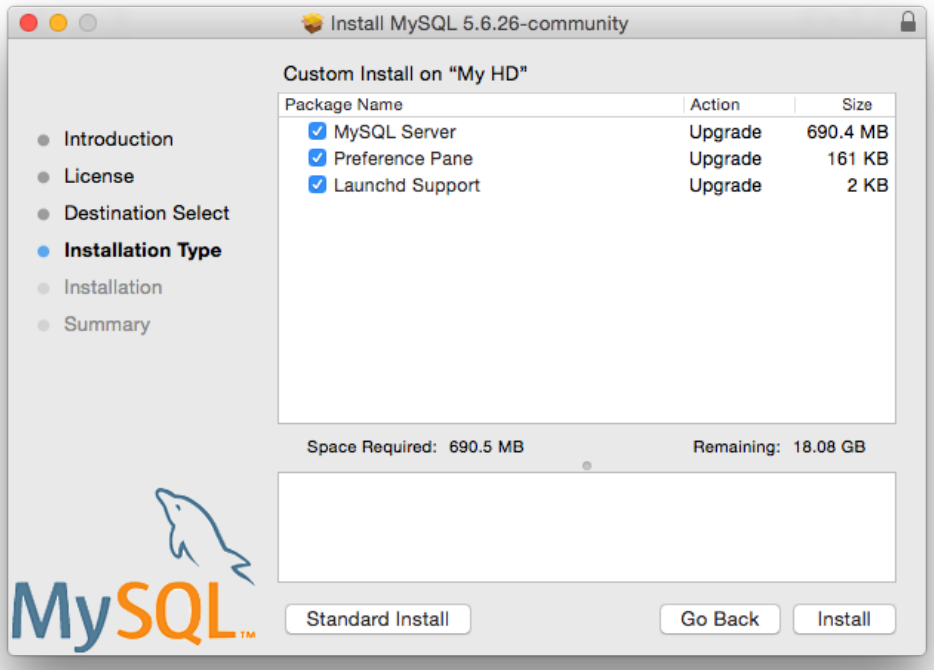
Figure 2.32 MySQL Package Installer: DMG Contents



2. Go through the process of installing the MySQL server, as described in the documentation at Section 2.4.2, “Installing MySQL on OS X Using Native Packages”.

3. Click **Customize** at the **Installation Type** step. The "Preference Pane" option is listed there and enabled by default; make sure it is not deselected.

Figure 2.33 MySQL Installer on OS X: Customize



4. Complete the MySQL server installation process.

Note

The MySQL preference pane only starts and stops MySQL installation installed from the MySQL package installation that have been installed in the default location.

Once the MySQL preference pane has been installed, you can control your MySQL server instance using the preference pane. To use the preference pane, open the **System Preferences...** from the Apple menu. Select the MySQL preference pane by clicking the MySQL logo within the **bottom** section of the preference panes list.

Figure 2.34 MySQL Preference Pane: Location

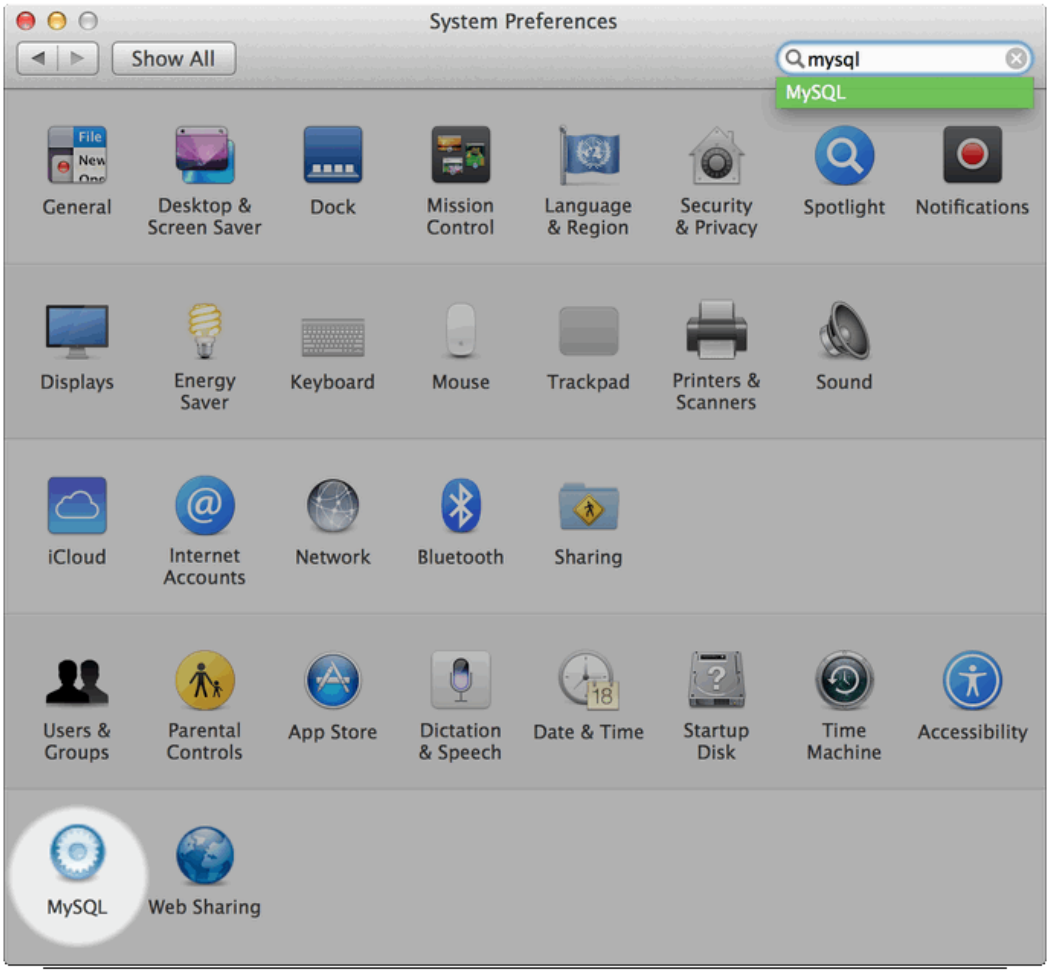
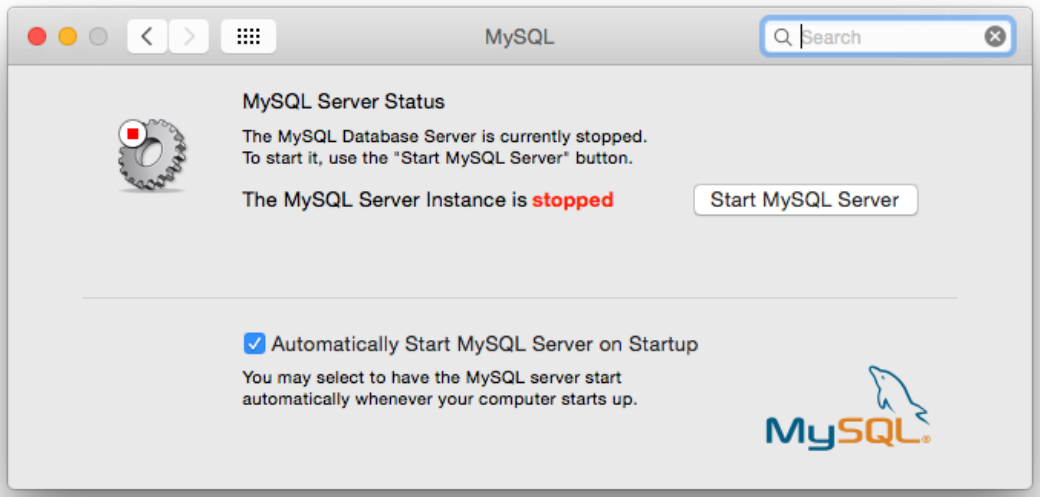


Figure 2.35 MySQL Preference Pane: Usage



The MySQL Preference Pane shows the current status of the MySQL server, showing **stopped** (in red) if the server is not running and **running** (in green) if the server has already been started. The preference pane also shows the current setting for whether the MySQL server has been set to start automatically.

- **To start the MySQL server using the preference pane:**

  Click **Start MySQL Server**. You may be prompted for the username and password of a user with administrator privileges to start the MySQL server.

- **To stop the MySQL server using the preference pane:**

  Click **Stop MySQL Server**. You may be prompted for the username and password of a user with administrator privileges to stop the MySQL server.

- **To automatically start the MySQL server when the system boots:**

  Check the check box next to **Automatically Start MySQL Server on Startup**.

- **To disable automatic MySQL server startup when the system boots:**

  Uncheck the check box next to **Automatically Start MySQL Server on Startup**.

You can close the `System Preferences...` window once you have completed your settings.

2.5 Installing MySQL on Linux

Linux supports a number of different solutions for installing MySQL. We recommend that you use one of the distributions from Oracle, for which several methods for installation are available:

- Installing with Yum using the MySQL Yum repository. For details, see Section 2.5.1, “Installing MySQL on Linux Using the MySQL Yum Repository”.
- Installing with APT using the MySQL APT Repository. For details, see Section 2.5.2, “Installing MySQL on Linux Using the MySQL APT Repository”.
- Installing with Zypper using the MySQL SLES Repository. For details, see Section 2.5.3, “Installing MySQL on Linux Using the MySQL SLES Repository”.

- Installing using a precompiled RPM package. For more information, see Section 2.5.4, “Installing MySQL on Linux Using RPM Packages from Oracle”.
- Installing using a precompiled Debian package. For more information, see Section 2.5.5, “Installing MySQL on Linux Using Debian Packages from Oracle”.
- Installing from a generic binary package in `.tar.gz` format. See Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries” for more information.
- Extracting and compiling MySQL from a source distribution. For detailed instructions, see Section 2.8, “Installing MySQL from Source”.

As an alternative, you can use the package manager on your system to automatically download and install MySQL with packages from the native software repositories of your Linux distribution. These native packages are often several versions behind the currently available release. You will also normally be unable to install development milestone releases (DMRs), as these are not usually made available in the native repositories. For more information on using the native package installers, see Section 2.5.6, “Installing MySQL on Linux from the Native Software Repositories”.

Note

For many Linux installations, you will want to set up MySQL to be started automatically when your machine starts. Many of the native package installations perform this operation for you, but for source, binary and RPM solutions you may need to set this up separately. The required script, `mysql.server`, can be found in the `support-files` directory under the MySQL installation directory or in a MySQL source tree. You can install it as `/etc/init.d/mysql` for automatic MySQL startup and shutdown. See Section 4.3.3, “`mysql.server` — MySQL Server Startup Script”.

2.5.1 Installing MySQL on Linux Using the MySQL Yum Repository

MySQL provides a Yum-style software repository for the following Linux distributions:

- EL6 and EL7-based platforms (for example, the corresponding versions of Red Hat Enterprise Linux, Oracle Linux, and CentOS)
- Fedora 23, 24, and 25

Currently, the MySQL Yum repository for the above-mentioned platforms provides RPM packages for installing the MySQL server, client, MySQL Workbench, MySQL Utilities, Connector/ODBC, and Connector/Python (not all packages are available for all the platforms; see Installing Additional MySQL Products and Components with Yum for details).

Before You Start

As a popular, open-source software, MySQL, in its original or re-packaged form, is widely installed on many systems from various sources, including different software download sites, software repositories, and so on. The following instructions assume that MySQL is not already installed on your system using a third-party-distributed RPM package; if that is not the case, follow the instructions given in Section 2.10.1.2, “Upgrading MySQL with the MySQL Yum Repository” or Replacing a Third-Party Distribution of MySQL Using the MySQL Yum Repository.

Steps for a Fresh Installation of MySQL

Follow the steps below to install the latest GA version of MySQL with the MySQL Yum repository:

Adding the MySQL Yum Repository

First, add the MySQL Yum repository to your system's repository list. This is a one-time operation, which can be performed by installing an RPM provided by MySQL. Follow these steps:

a. Go to the Download MySQL Yum Repository page (http://dev.mysql.com/downloads/repo/yum/) in the MySQL Developer Zone.

b. Select and download the release package for your platform.

c. Install the downloaded release package with the following command, replacing *platform-and-version-specific-package-name* with the name of the downloaded RPM package:

```
shell> sudo yum localinstall platform-and-version-specific-package-name.rpm
```

For an EL6-based system, the command is in the form of:

```
shell> sudo yum localinstall mysql57-community-release-el6-{version-number}.noarch.rpm
```

For an EL7-based system:

```
shell> sudo yum localinstall mysql57-community-release-el7-{version-number}.noarch.rpm
```

For Fedora 23:

```
shell> sudo dnf localinstall mysql57-community-release-fc23-{version-number}.noarch.rpm
```

For Fedora 24:

```
shell> sudo dnf localinstall mysql57-community-release-fc24-{version-number}.noarch.rpm
```

For Fedora 25:

```
shell> sudo dnf localinstall mysql57-community-release-fc25-{version-number}.noarch.rpm
```

The installation command adds the MySQL Yum repository to your system's repository list and downloads the GnuPG key to check the integrity of the software packages. See Section 2.1.3.2, “Signature Checking Using GnuPG” for details on GnuPG key checking.

You can check that the MySQL Yum repository has been successfully added by the following command (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> yum repolist enabled | grep "mysql.*-community.*"
```

Note

Once the MySQL Yum repository is enabled on your system, any system-wide update by the `yum update` command (or `dnf upgrade` for dnf-enabled systems) will upgrade MySQL packages on your system and also replace any native third-party packages, if Yum finds replacements for them in the MySQL Yum repository; see Section 2.10.1.2, “Upgrading MySQL with the MySQL Yum Repository” and, for a discussion on some possible effects of that on your system, see Upgrading the Shared Client Libraries.

2.4.3.3 Selecting a Release Series

When using the MySQL Yum repository, the latest GA series (currently MySQL 5.7) is selected for installation by default. If this is what you want, you can skip to the next step, Installing MySQL.

Within the MySQL Yum repository, different release series of the MySQL Community Server are hosted in different subrepositories. The subrepository for the latest GA series (currently MySQL

5.7) is enabled by default, and the subrepositories for all other series (for example, the MySQL 8.0 series) are disabled by default. Use this command to see all the subrepositories in the MySQL Yum repository, and see which of them are enabled or disabled (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> yum repolist all | grep mysql
```

To install the latest release from the latest GA series, no configuration is needed. To install the latest release from a specific series other than the latest GA series, disable the subrepository for the latest GA series and enable the subrepository for the specific series before running the installation command. If your platform supports `yum-config-manager`, you can do that by issuing these commands, which disable the subrepository for the 5.7 series and enable the one for the 8.0 series:

```
shell> sudo yum-config-manager --disable mysql57-community
shell> sudo yum-config-manager --enable mysql80-community
```

For dnf-enabled platforms:

```
shell> sudo dnf config-manager --disable mysql57-community
shell> sudo dnf config-manager --enable mysql80-community
```

Besides using `yum-config-manager` or the `dnf config-manager` command, you can also select a release series by editing manually the `/etc/yum.repos.d/mysql-community.repo` file. This is a typical entry for a release series' subrepository in the file:

```
[mysql57-community]
name=MySQL 5.7 Community Server
baseurl=http://repo.mysql.com/yum/mysql-5.7-community/el/6/$basearch/
enabled=1
gpgcheck=1
gpgkey=file:///etc/pki/rpm-gpg/RPM-GPG-KEY-mysql
```

Find the entry for the subrepository you want to configure, and edit the `enabled` option. Specify `enabled=0` to disable a subrepository, or `enabled=1` to enable a subrepository. For example, to install MySQL 8.0, make sure you have `enabled=0` for the above subrepository entry for MySQL 5.7, and have `enabled=1` for the entry for the 8.0 series:

```
# Enable to use MySQL 8.0
[mysql80-community]
name=MySQL 8.0 Community Server
baseurl=http://repo.mysql.com/yum/mysql-8.0-community/el/6/$basearch/
enabled=1
gpgcheck=1
gpgkey=file:///etc/pki/rpm-gpg/RPM-GPG-KEY-mysql
```

You should only enable subrepository for one release series at any time. When subrepositories for more than one release series are enabled, the latest series will be used by Yum.

Verify that the correct subrepositories have been enabled and disabled by running the following command and checking its output (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> yum repolist enabled | grep mysql
```

Installing MySQL

Install MySQL by the following command (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> sudo yum install mysql-community-server
```

This installs the package for MySQL server (`mysql-community-server`) and also packages for the components required to run the server, including packages for the client (`mysql-community-client`), the common error messages and character sets for client and server (`mysql-community-common`), and the shared client libraries (`mysql-community-libs`).

Starting the MySQL Server

Start the MySQL server with the following command:

```
shell> sudo service mysqld start
Starting mysqld:[ OK ]
```

You can check the status of the MySQL server with the following command:

```
shell> sudo service mysqld status
mysqld (pid 3066) is running.
```

At the initial start up of the server, the following happens, given that the data directory of the server is empty:

- The server is initialized.
- SSL certificate and key files are generated in the data directory.
- The validate_password plugin is installed and enabled.
- A superuser account `'root'@'localhost'` is created. A password for the superuser is set and stored in the error log file. To reveal it, use the following command:

```
shell> sudo grep 'temporary password' /var/log/mysqld.log
```

Change the root password as soon as possible by logging in with the generated, temporary password and set a custom password for the superuser account:

```
shell> mysql -uroot -p
```

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'MyNewPass4!';
```

Note

MySQL's validate_password plugin is installed by default. This will require that passwords contain at least one upper case letter, one lower case letter, one digit, and one special character, and that the total password length is at least 8 characters.

For more information on the postinstallation procedures, see Section 2.9, “Postinstallation Setup and Testing”.

Note

Compatibility Information for EL7-based platforms: The following RPM packages from the native software repositories of the platforms are incompatible with the package from the MySQL Yum repository that installs the MySQL server. Once you have installed MySQL using the MySQL Yum repository, you will not be able to install these packages (and vice versa).

- akonadi-mysql

Installing Additional MySQL Products and Components with Yum

You can use Yum to install and manage individual components of MySQL. Some of these components are hosted in sub-repositories of the MySQL Yum repository: for example, the MySQL Connectors are to be found in the MySQL Connectors Community sub-repository, and the MySQL Workbench in MySQL Tools Community. You can use the following command to list the packages for all the MySQL components available for your platform from the MySQL Yum repository (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> sudo yum --disablerepo=\* --enablerepo='mysql*-community*' list available
```

Install any packages of your choice with the following command, replacing *`package-name`* with name of the package (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> sudo yum install package-name
```

For example, to install MySQL Workbench on Fedora 24:

```
shell> sudo dnf install mysql-workbench-community
```

To install the shared client libraries (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
shell> sudo yum install mysql-community-libs
```

Updating MySQL with Yum

Besides installation, you can also perform updates for MySQL products and components using the MySQL Yum repository. See Section 2.10.1.2, “Upgrading MySQL with the MySQL Yum Repository” for details.

2.5.2 Installing MySQL on Linux Using the MySQL APT Repository

The MySQL APT repository provides `deb` packages for installing and managing the MySQL server, client, and other components on the following Linux platforms: :

- Debian 8.x (“jessie”)
- Ubuntu 14.04 LTS (“Trusty Tahr”)
- Ubuntu 16.04 (“Xenial Xerus”)

Instructions for using the MySQL APT Repository are available in A Quick Guide to Using the MySQL APT Repository.

2.5.3 Installing MySQL on Linux Using the MySQL SLES Repository

The MySQL SLES repository provides RPM packages for installing and managing the MySQL server, client, and other components on SUSE Enterprise Linux Server.

Instructions for using the MySQL SLES repository are available in A Quick Guide to Using the MySQL SLES Repository.

Note

The MySQL SLES repository is now in development release. We encourage you to try it and provide us with feedback. Please report any bugs or inconsistencies you observe to our Bugs Database.

2.5.4 Installing MySQL on Linux Using RPM Packages from Oracle

The recommended way to install MySQL on RPM-based Linux distributions is by using the RPM packages provided by Oracle. There are two sources for obtaining them, for the Community Edition of MySQL:

- From the MySQL software repositories:
  - The MySQL Yum repository (see Section 2.5.1, “Installing MySQL on Linux Using the MySQL Yum Repository” for details).
  - The MySQL SLES repository (see Section 2.5.3, “Installing MySQL on Linux Using the MySQL SLES Repository” for details).
- From the Download MySQL Community Server page in the MySQL Developer Zone.

Note

RPM distributions of MySQL are also provided by other vendors. Be aware that they may differ from those built by Oracle in features, capabilities, and conventions (including communication setup), and that the installation instructions in this manual do not necessarily apply to them. The vendor's instructions should be consulted instead.

RPM packages for MySQL are listed in the following tables:

Table 2.7 RPM Packages for MySQL Community Edition

Package Name	Summary
mysql-community-server	Database server and related tools
mysql-community-client	MySQL client applications and tools
mysql-community-common	Common files for server and client libraries
mysql-community-devel	Development header files and libraries for MySQL database client applications
mysql-community-libs	Shared libraries for MySQL database client applications
mysql-community-libs-compat	Shared compatibility libraries for previous MySQL installations
mysql-community-test	Test suite for the MySQL server

Table 2.8 RPM Packages for the MySQL Enterprise Edition

Package Name	Summary
mysql-commercial-server	Database server and related tools
mysql-commercial-client	MySQL client applications and tools
mysql-commercial-common	Common files for server and client libraries
mysql-commercial-devel	Development header files and libraries for MySQL database client applications
mysql-commercial-libs	Shared libraries for MySQL database client applications
mysql-commercial-libs-compat	Shared compatibility libraries for previous MySQL installations
mysql-commercial-test	Test suite for the MySQL server

The full names for the RPMs have the following syntax:

```
packagename-version-distribution-arch.rpm
```

The *distribution* and *arch* values indicate the Linux distribution and the processor type for which the package was built. See the table below for lists of the distribution identifiers:

Table 2.9 MySQL Linux RPM Package Distribution Identifiers

distribution Value	Intended Use
el6, el7	Red Hat Enterprise Linux/Oracle Linux/CentOS 6 or 7
fc23, fc24, fc25	Fedora 23, 24, or 25
sles12	SUSE Linux Enterprise Server 12

To see all files in an RPM package (for example, mysql-community-server), use the following command:

```
shell> rpm -qpl mysql-community-server-version-distribution-arch.rpm
```

The discussion in the rest of this section applies only to an installation process using the RPM packages directly downloaded from Oracle, instead of through a MySQL repository.

Dependency relationships exist among some of the packages. If you plan to install many of the packages, you may wish to download the RPM bundle tar file instead, which contains all the RPM packages listed above, so that you need not download them separately.

In most cases, you need to install the mysql-community-server, mysql-community-client, mysql-community-libs, mysql-community-common, and mysql-community-libs-compat packages to get a functional, standard MySQL installation. To perform such a standard, minimal installation, go to the folder that contains all those packages (and, preferably, no other RPM packages with similar names), and issue the following command:

```
shell> sudo yum install mysql-community-{server,client,common,libs}-*
```

Replace yum with zypper for SLES systems, and with dnf for dnf-enabled systems (like Fedora 23 and later).

While it is much preferable to use a high-level package management tool like yum to install the packages, users who prefer direct rpm commands can replace the yum install command with the rpm -Uvh command; however, using rpm -Uvh instead makes the installation process more prone to failure, due to potential dependency issues the installation process might run into.

To install only the client programs, you can skip mysql-community-server in your list of packages to install; issue the following command:

```
shell> sudo yum install mysql-community-{client,common,libs}-*
```

Replace yum with zypper for SLES systems, and with dnf for dnf-enabled systems (like Fedora 23 and later).

A standard installation of MySQL using the RPM packages result in files and resources created under the system directories, shown in the following table.

Table 2.10 MySQL Installation Layout for Linux RPM Packages from the MySQL Developer Zone

Files or Resources	Location
Client programs and scripts	/usr/bin
mysqld server	/usr/sbin
Configuration file	/etc/my.cnf

Files or Resources	Location
Data directory	`/var/lib/mysql`
Error log file	For RHEL, Oracle Linux, CentOS or Fedora platforms: `/var/log/mysqld.log` For SLES: `/var/log/mysql/mysqld.log`
Value of `secure_file_priv`	`/var/lib/mysql-files`
System V init script	For RHEL, Oracle Linux, CentOS or Fedora platforms: `/etc/init.d/mysqld` For SLES: `/etc/init.d/mysql`
Systemd service	For RHEL, Oracle Linux, CentOS or Fedora platforms: `mysqld` For SLES: `mysql`
Pid file	`/var/run/mysql/mysqld.pid`
Socket	`/var/lib/mysql/mysql.sock`
Keyring directory	`/var/lib/mysql-keyring`
Unix manual pages	`/usr/share/man`
Include (header) files	`/usr/include/mysql`
Libraries	`/usr/lib/mysql`
Miscellaneous support files (for example, error messages, and character set files)	`/usr/share/mysql`

The installation also creates a user named `mysql` and a group named `mysql` on the system.

Note

Installation of previous versions of MySQL using older packages might have created a configuration file named `/usr/my.cnf`. It is highly recommended that you examine the contents of the file and migrate the desired settings inside to the file `/etc/my.cnf` file, then remove `/usr/my.cnf`.

MySQL is NOT automatically started at the end of the installation process. For Red Hat Enterprise Linux, Oracle Linux, CentOS, and Fedora systems, use the following command to start MySQL:

```
shell> sudo service mysqld start
```

For SLES systems, the command is the same, but the service name is different:

```
shell> sudo service mysql start
```

If the operating system is systemd enabled, standard `service` commands such as `stop`, `start`, `status` and `restart` should be used to manage the MySQL server service. The `mysqld` service is enabled by default, and it starts at system reboot. Notice that certain things might work differently on systemd platforms: for example, changing the location of the data directory might cause issues. See Section 2.5.9, “Managing MySQL Server with systemd” for additional information.

At the initial start up of the server, the following happens, given that the data directory of the server is empty:

- The server is initialized.
- An SSL certificate and key files are generated in the data directory.
- The validate_password plugin is installed and enabled.

- A superuser account `'root'@'localhost'` is created. A password for the superuser is set and stored in the error log file. To reveal it, use the following command for RHEL, Oracle Linux, CentOS, and Fedora systems:

```
shell> sudo grep 'temporary password' /var/log/mysqld.log
```

  Use the following command for SLES systems:

```
shell> sudo grep 'temporary password' /var/log/mysql/mysqld.log
```

  The next step is to log in with the generated, temporary password and set a custom password for the superuser account:

```
shell> mysql -uroot -p
```

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'MyNewPass4!';
```

Note

MySQL's validate_password plugin is installed by default. This will require that passwords contain at least one upper case letter, one lower case letter, one digit, and one special character, and that the total password length is at least 8 characters.

If something goes wrong during installation, you might find debug information in the error log file `/var/log/mysqld.log`.

For some Linux distributions, it might be necessary to increase the limit on number of file descriptors available to `mysqld`. See Section B.5.2.17, “File Not Found and Similar Errors”

Installing Client Libraries from Multiple MySQL Versions. It is possible to install multiple client library versions, such as for the case that you want to maintain compatibility with older applications linked against previous libraries. To install an older client library, use the `--oldpackage` option with `rpm`. For example, to install `mysql-community-libs-5.5` on an EL6 system that has `libmysqlclient.21` from MySQL 8.0, use a command like this:

```
shell> rpm --oldpackage -ivh mysql-community-libs-5.5.50-2.el6.x86_64.rpm
```

Debug Package. A special variant of MySQL Server compiled with the debug package has been included in the server RPM packages. It performs debugging and memory allocation checks and produces a trace file when the server is running. To use that debug version, start MySQL with `/usr/sbin/mysqld-debug`, instead of starting it as a service or with `/usr/sbin/mysqld`. See Section 28.5.3, “The DBUG Package” for the debug options you can use.

Rebuilding RPMs from source SRPMs. Source code SRPM packages for MySQL are available for download. They can be used as-is to rebuild the MySQL RPMs with the standard `rpmbuild` tool chain.

2.5.5 Installing MySQL on Linux Using Debian Packages from Oracle

Oracle provides Debian packages for installing MySQL on Debian or Debian-like Linux systems. The packages are available through two different channels:

- The MySQL APT Repository. This is the preferred method for installing MySQL on Debian-like systems, as it provides a simple and convenient way to install and update MySQL products. For details, see Section 2.5.2, “Installing MySQL on Linux Using the MySQL APT Repository”.
- The MySQL Developer Zone's Download Area. For details, see Section 2.1.2, “How to Get MySQL”. The following are some information on the Debian packages available there and the instructions for installing them:

- Various Debian packages are provided in the MySQL Developer Zone for installing different components of MySQL on different Debian or Ubuntu platforms (currently, Debian 8, and Ubuntu 14 and 16 are supported). The preferred method is to use the tarball bundle, which contains the packages needed for a basic setup of MySQL. The tarball bundles have names in the format of `mysql-server_MVER-DVER_CPU.deb-bundle.tar`. *MVER* is the MySQL version and *DVER* is the Linux distribution version. The *CPU* value indicates the processor type or family for which the package is built, as shown in the following table:

  Table 2.11 MySQL Debian and Ubuntu Installation Packages CPU Identifiers

CPU Value	Intended Processor Type or Family
`i386`	Pentium processor or better, 32 bit
`amd64`	64-bit x86 processor

- After downloading the tarball, unpack it with the following command:

```
shell> tar -xvf mysql-server_MVER-DVER_CPU.deb-bundle.tar
```

- You may need to install the `libaio` library if it is not already present on your system:

```
shell> sudo apt-get install libaio1
```

- Preconfigure the MySQL server package with the following command:

```
shell> sudo dpkg-preconfigure mysql-community-server_*.deb
```

  You will be asked to provide a password for the root user for your MySQL installation. You might also be asked other questions regarding the installation.

  Important

  Make sure you remember the root password you set. Users who want to set a password later can leave the **password** field blank in the dialogue box and just press **OK**; in that case, root access to the server is authenticated using the MySQL Socket Peer-Credential Authentication Plugin for connections using a Unix socket file. You can set the root password later using `mysql_secure_installation`.

- For a basic installation of the MySQL server, install the database common files package, the client package, the client metapackage, the server package, and the server metapackage (in that order); you can do that with a single command:

```
shell> sudo dpkg -i mysql-{common,community-client,client,community-server,server}_*.deb
```

  There are also packages with `server-core` and `client-core` in the package names. These contain binaries only and are installed automatically by the standard packages. Installing them by themselves will not result in a functioning MySQL setup.

  If you are being warned of unmet dependencies by `dpkg`, you can fix them using `apt-get`:

```
sudo apt-get -f install
```

  Here are where the files are installed on the system:

  - All configuration files (like `my.cnf`) are under `/etc/mysql`
  - All binaries, libraries, headers, etc., are under `/usr/bin` and `/usr/sbin`

- The data directory is under `/var/lib/mysql`

Note

Debian distributions of MySQL are also provided by other vendors. Be aware that they may differ from those built by Oracle in features, capabilities, and conventions (including communication setup), and that the instructions in this manual do not necessarily apply to installing them. The vendor's instructions should be consulted instead.

2.5.6 Installing MySQL on Linux from the Native Software Repositories

Many Linux distributions include a version of the MySQL server, client tools, and development components in their native software repositories and can be installed with the platforms' standard package management systems. This section provides basic instructions for installing MySQL using those package management systems.

Important

Native packages are often several versions behind the currently available release. You will also normally be unable to install development milestone releases (DMRs), as these are not usually made available in the native repositories. Before proceeding, we recommend that you check out the other installation options described in Section 2.5, “Installing MySQL on Linux”.

Distribution specific instructions are shown below:

- **Red Hat Linux, Fedora, CentOS**

Note

For a number of Linux distributions, you can install MySQL using the MySQL Yum repository instead of the platform's native software repository. See Section 2.5.1, “Installing MySQL on Linux Using the MySQL Yum Repository” for details.

For Red Hat and similar distributions, the MySQL distribution is divided into a number of separate packages, `mysql` for the client tools, `mysql-server` for the server and associated tools, and `mysql-libs` for the libraries. The libraries are required if you want to provide connectivity from different languages and environments such as Perl, Python and others.

To install, use the `yum` command to specify the packages that you want to install. For example:

```
root-shell> yum install mysql mysql-server mysql-libs mysql-server
Loaded plugins: presto, refresh-packagekit
Setting up Install Process
Resolving Dependencies
--> Running transaction check
---> Package mysql.x86_64 0:5.1.48-2.fc13 set to be updated
---> Package mysql-libs.x86_64 0:5.1.48-2.fc13 set to be updated
---> Package mysql-server.x86_64 0:5.1.48-2.fc13 set to be updated
--> Processing Dependency: perl-DBD-MySQL for package: mysql-server-5.1.48-2.fc13.x86_64
--> Running transaction check
---> Package perl-DBD-MySQL.x86_64 0:4.017-1.fc13 set to be updated
--> Finished Dependency Resolution

Dependencies Resolved

================================================================================
 Package               Arch          Version               Repository      Size
================================================================================
Installing:
 mysql                 x86_64        5.1.48-2.fc13         updates        889 k
```

```
 mysql-libs            x86_64        5.1.48-2.fc13         updates        1.2 M
 mysql-server          x86_64        5.1.48-2.fc13         updates        8.1 M
Installing for dependencies:
 perl-DBD-MySQL        x86_64        4.017-1.fc13          updates        136 k

Transaction Summary
================================================================================
Install       4 Package(s)
Upgrade       0 Package(s)

Total download size: 10 M
Installed size: 30 M
Is this ok [y/N]: y
Downloading Packages:
Setting up and reading Presto delta metadata
Processing delta metadata
Package(s) data still to download: 10 M
(1/4): mysql-5.1.48-2.fc13.x86_64.rpm                     | 889 kB     00:04
(2/4): mysql-libs-5.1.48-2.fc13.x86_64.rpm                | 1.2 MB     00:06
(3/4): mysql-server-5.1.48-2.fc13.x86_64.rpm              | 8.1 MB     00:40
(4/4): perl-DBD-MySQL-4.017-1.fc13.x86_64.rpm             | 136 kB     00:00
--------------------------------------------------------------------------------
Total                                          201 kB/s |  10 MB     00:52
Running rpm_check_debug
Running Transaction Test
Transaction Test Succeeded
Running Transaction
  Installing     : mysql-libs-5.1.48-2.fc13.x86_64                          1/4
  Installing     : mysql-5.1.48-2.fc13.x86_64                               2/4
  Installing     : perl-DBD-MySQL-4.017-1.fc13.x86_64                       3/4
  Installing     : mysql-server-5.1.48-2.fc13.x86_64                        4/4

Installed:
  mysql.x86_64 0:5.1.48-2.fc13            mysql-libs.x86_64 0:5.1.48-2.fc13
  mysql-server.x86_64 0:5.1.48-2.fc13

Dependency Installed:
  perl-DBD-MySQL.x86_64 0:4.017-1.fc13

Complete!
```

MySQL and the MySQL server should now be installed. A sample configuration file is installed into `/etc/my.cnf`. An init script, to start and stop the server, will have been installed into `/etc/init.d/mysqld`. To start the MySQL server use `service`:

```
root-shell> service mysqld start
```

To enable the server to be started and stopped automatically during boot, use `chkconfig`:

```
root-shell> chkconfig --levels 235 mysqld on
```

Which enables the MySQL server to be started (and stopped) automatically at the specified the run levels.

The database tables will have been automatically created for you, if they do not already exist. You should, however, run `mysql_secure_installation` to set the root passwords on your server.

- **Debian, Ubuntu, Kubuntu**

Note

For Debian 8 and Ubuntu 14 and 16, MySQL can be installed using the MySQL APT Repository instead of the platform's native software repository. See Section 2.5.2, “Installing MySQL on Linux Using the MySQL APT Repository” for details.

On Debian and related distributions, there are two packages for MySQL in their software repositories, `mysql-client` and `mysql-server`, for the client and server components respectively. You should specify an explicit version, for example `mysql-client-5.1`, to ensure that you install the version of MySQL that you want.

To download and install, including any dependencies, use the `apt-get` command, specifying the packages that you want to install.

Note

Before installing, make sure that you update your `apt-get` index files to ensure you are downloading the latest available version.

A sample installation of the MySQL packages might look like this (some sections trimmed for clarity):

```
root-shell> apt-get install mysql-client-5.1 mysql-server-5.1
Reading package lists... Done
Building dependency tree
Reading state information... Done
The following packages were automatically installed and are no longer required:
  linux-headers-2.6.28-11 linux-headers-2.6.28-11-generic
Use 'apt-get autoremove' to remove them.
The following extra packages will be installed:
  bsd-mailx libdbd-mysql-perl libdbi-perl libhtml-template-perl
  libmysqlclient15off libmysqlclient16 libnet-daemon-perl libplrpc-perl mailx
  mysql-common postfix
Suggested packages:
  dbishell libipc-sharedcache-perl tinyca procmail postfix-mysql postfix-pgsql
  postfix-ldap postfix-pcre sasl2-bin resolvconf postfix-cdb
The following NEW packages will be installed
  bsd-mailx libdbd-mysql-perl libdbi-perl libhtml-template-perl
  libmysqlclient15off libmysqlclient16 libnet-daemon-perl libplrpc-perl mailx
  mysql-client-5.1 mysql-common mysql-server-5.1 postfix
0 upgraded, 13 newly installed, 0 to remove and 182 not upgraded.
Need to get 1907kB/25.3MB of archives.
After this operation, 59.5MB of additional disk space will be used.
Do you want to continue [Y/n]? Y
Get: 1 http://gb.archive.ubuntu.com jaunty-updates/main mysql-common 5.1.30really5.0.75-0ubuntu10.5 [
Get: 2 http://gb.archive.ubuntu.com jaunty-updates/main libmysqlclient15off 5.1.30really5.0.75-0ubunt
Fetched 1907kB in 9s (205kB/s)
Preconfiguring packages ...
Selecting previously deselected package mysql-common.
(Reading database ... 121260 files and directories currently installed.)
...
Processing 1 added doc-base file(s)...
Registering documents with scrollkeeper...
Setting up libnet-daemon-perl (0.43-1) ...
Setting up libplrpc-perl (0.2020-1) ...
Setting up libdbi-perl (1.607-1) ...
Setting up libmysqlclient15off (5.1.30really5.0.75-0ubuntu10.5) ...

Setting up libdbd-mysql-perl (4.008-1) ...
Setting up libmysqlclient16 (5.1.31-1ubuntu2) ...

Setting up mysql-client-5.1 (5.1.31-1ubuntu2) ...

Setting up mysql-server-5.1 (5.1.31-1ubuntu2) ...
 * Stopping MySQL database server mysqld
   ...done.
2013-09-24T13:03:09.048353Z 0 [Note] InnoDB: 8.0.3 started; log sequence number 1566036
2013-09-24T13:03:10.057269Z 0 [Note] InnoDB: Starting shutdown...
2013-09-24T13:03:10.857032Z 0 [Note] InnoDB: Shutdown completed; log sequence number 1566036
 * Starting MySQL database server mysqld
   ...done.
 * Checking for corrupt, not cleanly closed and upgrade needing tables.
...
Processing triggers for libc6 ...
```

```
ldconfig deferred processing now taking place
```

Note

The `apt-get` command will install a number of packages, including the MySQL server, in order to provide the typical tools and application environment. This can mean that you install a large number of packages in addition to the main MySQL package.

During installation, the initial database will be created, and you will be prompted for the MySQL root password (and confirmation). A configuration file will have been created in `/etc/mysql/my.cnf`. An init script will have been created in `/etc/init.d/mysql`.

The server will already be started. You can manually start and stop the server using:

```
root-shell> service mysql [start|stop]
```

The service will automatically be added to the 2, 3 and 4 run levels, with stop scripts in the single, shutdown and restart levels.

2.5.7 Installing MySQL on Linux with docker

The docker deployment framework supports easy installation and configuration of MySQL servers. For instructions, see https://hub.docker.com/r/mysql/mysql-server/. This page also provides extensive documentation about using MySQL under docker.

2.5.8 Installing MySQL on Linux with juju

The juju deployment framework supports easy installation and configuration of MySQL servers. For instructions, see https://jujucharms.com/mysql/.

2.5.9 Managing MySQL Server with systemd

If you install MySQL using an RPM or Debian package on the following Linux platforms, server startup and shutdown is managed by systemd:

- RPM package platforms:
  - Red Hat Enterprise Linux 7; Oracle Linux 7; CentOS 7
  - SUSE Linux Enterprise Server 12
  - Fedora 24 and 25
- Debian package platforms:
  - Debian 8 or higher
  - Ubuntu 16 or higher

If you install MySQL from a source distribution on a platform that uses systemd, obtain systemd support for MySQL by configuring the distribution using the `-DWITH_SYSTEMD=1 CMake` option. See Section 2.8.4, “MySQL Source-Configuration Options”.

The following discussion covers these topics:

- Overview of systemd
- Configuring systemd for MySQL
- Configuring Multiple MySQL Instances Using systemd

- Migrating from mysqld_safe to systemd

Note

On platforms for which systemd support for MySQL is installed, scripts such as `mysqld_safe` and the System V initialization script are unnecessary and are not installed. For example, `mysqld_safe` can handle server restarts, but systemd provides the same capability, and does so in a manner consistent with management of other services rather than by using an application-specific program.

Because systemd has the capability of managing multiple MySQL instances on platforms for which systemd support for MySQL is installed, `mysqld_multi` and `mysqld_multi.server` are unnecessary and are not installed.

Overview of systemd

systemd provides automatic MySQL server startup and shutdown. It also enables manual server management using the `systemctl` command. For example:

```
systemctl {start|stop|restart|status} mysqld
```

Alternatively, use the `service` command (with the arguments reversed), which is compatible with System V systems:

```
service mysqld {start|stop|restart|status}
```

Note

For the `systemctl` or `service` commands, if the MySQL service name is not `mysqld`, use the appropriate name. For example, use `mysql` rather than `mysqld` on Debian-based and SLES systems.

Support for systemd includes these files:

- `mysqld.service` (RPM platforms), `mysql.service` (Debian platforms): systemd service unit configuration file, with details about the MySQL service.
- `mysqld@.service` (RPM platforms), `mysql@.service` (Debian platforms): Like `mysqld.service` or `mysql.service`, but used for managing multiple MySQL instances.
- `mysqld.tmpfiles.d`: File containing information to support the `tmpfiles` feature. This file is installed under the name `mysql.conf`.
- `mysqld_pre_systemd` (RPM platforms), `mysql-system-start` (Debian platforms): Support script for the unit file. This script assists in creating the error log file only if the log location matches a pattern (`/var/log/mysql*.log` for RPM platforms, `/var/log/mysql/*.log` for Debian platforms). In other cases, the error log directory must be writable or the error log must be present and writable for the user running the `mysqld` process.

Configuring systemd for MySQL

To add or change systemd options for MySQL, these methods are available:

- Use a localized systemd configuration file.
- Arrange for systemd to set environment variables for the MySQL server process.
- Set the `MYSQLD_OPTS` systemd variable.

To use a localized systemd configuration file, create the `/etc/systemd/system/mysqld.service.d` directory if it does not exist. In that directory, create a file that contains a `[Service]` section listing the desired settings. For example:

```
[Service]
LimitNOFILE=max_open_files
PIDFile=/path/to/pid/file
Nice=nice_level
LimitCore=core_file_limit
Environment="LD_PRELOAD=/path/to/malloc/library"
Environment="TZ=time_zone_setting"
```

The discussion here uses `override.conf` as the name of this file. Newer versions of systemd support the following command, which opens an editor and permits you to edit the file:

```
systemctl edit mysqld  # RPM platforms
systemctl edit mysql   # Debian platforms
```

Whenever you create or change `override.conf`, reload the systemd configuration, then tell systemd to restart the MySQL service:

```
systemctl daemon-reload
systemctl restart mysqld  # RPM platforms
systemctl restart mysql   # Debian platforms
```

With systemd, the `override.conf` configuration method must be used for certain parameters, rather than settings in a `[mysqld]` or `[mysqld_safe]` group in a MySQL option file:

- For some parameters, `override.conf` must be used because systemd itself must know their values and it cannot read MySQL option files to get them.
- Parameters that specify values otherwise settable only using options known to `mysqld_safe` must be specified using systemd because there is no corresponding `mysqld` parameter.

For additional information about using systemd rather than `mysqld_safe`, see Migrating from mysqld_safe to systemd.

You can set the following parameters in `override.conf`:

- To specify the process ID file, use `override.conf` and change both `PIDFile` and `ExecStart` to name the PID file path name. Any setting of the process ID file in MySQL option files is ignored. To modify `ExecStart`, it must first be cleared. For example:

  ```
  [Service]
  PIDFile=/var/run/mysqld/mysqld-custom.pid
  ExecStart=
  ExecStart=/usr/sbin/mysqld --daemonize --pid-file=/var/run/mysqld/mysqld-custom.pid $MYSQLD_OPTS
  ```

- To set the number of file descriptors available to the MySQL server, use `LimitNOFILE` in `override.conf` rather than the `--open-files-limit` option for `mysqld` or `mysqld_safe`.
- To set the maximum core file size, use `LimitCore` in `override.conf` rather than the `--core-file-size` option for `mysqld_safe`.
- To set the scheduling priority for the MySQL server, use `Nice` in `override.conf` rather than the `--nice` option for `mysqld_safe`.

Some MySQL parameters are configured using environment variables:

- `LD_PRELOAD`: Set this variable if the MySQL server should use a specific memory-allocation library.
- `TZ`: Set this variable to specify the default time zone for the server.

There are multiple ways to specify environment variable values for use by the MySQL server process managed by systemd:

- Use `Environment` lines in the `override.conf` file. For the syntax, see the example in the preceding discussion that describes how to use this file.
- Specify the values in the `/etc/sysconfig/mysql` file (create the file if it does not exist). Assign values using the following syntax:

```
LD_PRELOAD=/path/to/malloc/library
TZ=time_zone_setting
```

  After modifying `/etc/sysconfig/mysql`, restart the server to make the changes effective:

```
systemctl restart mysqld  # RPM platforms
systemctl restart mysql   # Debian platforms
```

To specify options for `mysqld` without modifying systemd configuration files directly, set or unset the `MYSQLD_OPTS` systemd variable. For example:

```
systemctl set-environment MYSQLD_OPTS="--general_log=1"
systemctl unset-environment MYSQLD_OPTS
```

`MYSQLD_OPTS` can also be set in the `/etc/sysconfig/mysql` file.

After modifying the systemd environment, restart the server to make the changes effective:

```
systemctl restart mysqld  # RPM platforms
systemctl restart mysql   # Debian platforms
```

Configuring Multiple MySQL Instances Using systemd

This section describes how to configure systemd for multiple instances of MySQL.

Note

Because systemd has the capability of managing multiple MySQL instances on platforms for which systemd support is installed, `mysqld_multi` and `mysqld_multi.server` are unnecessary and are not installed.

To use multiple-instance capability, modify the `my.cnf` option file to include configuration of key options for each instance. These file locations are typical:

- `/etc/my.cnf` or `/etc/mysql/my.cnf` (RPM platforms)
- `/etc/mysql/mysql.conf.d/mysqld.cnf` (Debian platforms)

For example, to manage two instances named `replica01` and `replica02`, add something like this to the option file:

RPM platforms:

```
[mysqld@replica01]
datadir=/var/lib/mysql-replica01
socket=/var/lib/mysql-replica01/mysql.sock
port=3307
log-error=/var/log/mysqld-replica01.log

[mysqld@replica02]
datadir=/var/lib/mysql-replica02
```

```
socket=/var/lib/mysql-replica02/mysql.sock
port=3308
log-error=/var/log/mysqld-replica02.log
```

Debian platforms:

```
[mysqld@replica01]
datadir=/var/lib/mysql-replica01
socket=/var/lib/mysql-replica01/mysql.sock
port=3307
log-error=/var/log/mysql/replica01.log

[mysqld@replica02]
datadir=/var/lib/mysql-replica02
socket=/var/lib/mysql-replica02/mysql.sock
port=3308
log-error=/var/log/mysql/replica02.log
```

The replica names shown here use `@` as the delimiter because that is the only delimiter supported by systemd.

Instances then are managed by normal systemd commands, such as:

```
systemctl start mysqld@replica01
systemctl start mysqld@replica02
```

To enable instances to run at boot time, do this:

```
systemctl enable mysqld@replica01
systemctl enable mysqld@replica02
```

Use of wildcards is also supported. For example, this command displays the status of all replica instances:

```
systemctl status 'mysqld@replica*'
```

For management of multiple MySQL instances on the same machine, systemd automatically uses a different unit file:

- `mysqld@.service` rather than `mysqld.service` (RPM platforms)
- `mysql@.service` rather than `mysql.service` (Debian platforms)

In the unit file, `%I` and `%i` reference the parameter passed in after the `@` marker and are used to manage the specific instance. For a command such as this:

```
systemctl start mysqld@replica01
```

systemd starts the server using a command such as this:

```
mysqld --defaults-group-suffix=@%I ...
```

The result is that the `[server]`, `[mysqld]`, and `[mysqld@replica01]` option groups are read and used for that instance of the service.

Note

On Debian platforms, AppArmor prevents the server from reading or writing `/var/lib/mysql-replica*`, or anything other than the default locations. To

address this, you must customize or disable the profile in `/etc/apparmor.d/user.sbin.mysqld`.

Note

On Debian platforms, the packaging scripts for MySQL uninstallation cannot currently handle `mysqld@` instances. Before removing or upgrading the package, you must stop any extra instances manually first.

Migrating from mysqld_safe to systemd

Because `mysqld_safe` is not installed on platforms that use systemd to manage MySQL, options previously specified for that program (for example, in an `[mysqld_safe]` option group) must be specified another way:

- Some `mysqld_safe` options are also understood by `mysqld` and can be moved from the `[mysqld_safe]` option group to the `[mysqld]` group. This does *not* include `--pid-file`, `--open-files-limit`, or `--nice`. To specify those options, use the `override.conf` systemd file, described previously.
- For some `mysqld_safe` options, there are similar `mysqld` options. For example, the `mysqld_safe` option for enabling `syslog` logging is `--syslog`. For `mysqld`, enable the `log_syslog` system variable instead. For details, see Section 5.4.2, “The Error Log”.
- `mysqld_safe` options not understood by `mysqld` can be specified in `override.conf` or environment variables. For example, with `mysqld_safe`, if the server should use a specific memory allocation library, this is specified using the `--malloc-lib` option. For installations that manage the server with systemd, arrange to set the `LD_PRELOAD` environment variable instead, as described previously.

2.6 Installing MySQL on Solaris

Note

MySQL 8.0 supports Solaris 11 (Update 3+ and later).

MySQL on Solaris is available in a number of different formats.

- For information on installing using the native Solaris PKG format, see Section 2.6.1, “Installing MySQL on Solaris Using a Solaris PKG”.
- To use a standard `tar` binary installation, use the notes provided in Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”. Check the notes and hints at the end of this section for Solaris specific notes that you may need before or after installation.

Important

The installation packages have a dependency on the Oracle Developer Studio 12.5 Runtime Libraries, which must be installed before you run the MySQL installation package. See the download options for Oracle Developer Studio here. The installation package enables you to install the runtime libraries only instead of the full Oracle Developer Studio; see instructions in Installing Only the Runtime Libraries on Oracle Solaris 11.

To obtain a binary MySQL distribution for Solaris in tarball or PKG format, http://dev.mysql.com/downloads/mysql/8.0.html.

Additional notes to be aware of when installing and using MySQL on Solaris:

- If you want to use MySQL with the `mysql` user and group, use the `groupadd` and `useradd` commands:

```
groupadd mysql
useradd -g mysql -s /bin/false mysql
```

- If you install MySQL using a binary tarball distribution on Solaris, because the Solaris `tar` cannot handle long file names, use GNU `tar` (`gtar`) to unpack the distribution. If you do not have GNU `tar` on your system, install it with the following command:

```
pkg install archiver/gnu-tar
```

- You should mount any file systems on which you intend to store `InnoDB` files with the `forcedirectio` option. (By default mounting is done without this option.) Failing to do so will cause a significant drop in performance when using the `InnoDB` storage engine on this platform.
- If you would like MySQL to start automatically, you can copy `support-files/mysql.server` to `/etc/init.d` and create a symbolic link to it named `/etc/rc3.d/S99mysql.server`.
- If too many processes try to connect very rapidly to `mysqld`, you should see this error in the MySQL log:

```
Error in accept: Protocol error
```

  You might try starting the server with the `--back_log=50` option as a workaround for this.

- To configure the generation of core files on Solaris you should use the `coreadm` command. Because of the security implications of generating a core on a `setuid()` application, by default, Solaris does not support core files on `setuid()` programs. However, you can modify this behavior using `coreadm`. If you enable `setuid()` core files for the current user, they will be generated using the mode 600 and owned by the superuser.

2.6.1 Installing MySQL on Solaris Using a Solaris PKG

You can install MySQL on Solaris using a binary package of the native Solaris PKG format instead of the binary tarball distribution.

Important

The installation package has a dependency on the Oracle Developer Studio 12.5 Runtime Libraries, which must be installed before you run the MySQL installation package. See the download options for Oracle Developer Studio here. The installation package enables you to install the runtime libraries only instead of the full Oracle Developer Studio; see instructions in Installing Only the Runtime Libraries on Oracle Solaris 11.

To use this package, download the corresponding `mysql-VERSION-solaris11-PLATFORM.pkg.gz` file, then uncompress it. For example:

```
shell> gunzip mysql-8.0.3-solaris11-x86_64.pkg.gz
```

To install a new package, use `pkgadd` and follow the onscreen prompts. You must have root privileges to perform this operation:

```
shell> pkgadd -d mysql-8.0.3-solaris11-x86_64.pkg

The following packages are available:
  1  mysql     MySQL Community Server (GPL)
               (i86pc) 8.0.3

Select package(s) you wish to process (or 'all' to process
```

```
all packages). (default: all) [?,??,q]:
```

The PKG installer installs all of the files and tools needed, and then initializes your database if one does not exist. To complete the installation, you should set the root password for MySQL as provided in the instructions at the end of the installation. Alternatively, you can run the `mysql_secure_installation` script that comes with the installation.

By default, the PKG package installs MySQL under the root path `/opt/mysql`. You can change only the installation root path when using `pkgadd`, which can be used to install MySQL in a different Solaris zone. If you need to install in a specific directory, use a binary `tar` file distribution.

The `pkg` installer copies a suitable startup script for MySQL into `/etc/init.d/mysql`. To enable MySQL to startup and shutdown automatically, you should create a link between this file and the init script directories. For example, to ensure safe startup and shutdown of MySQL you could use the following commands to add the right links:

```
shell> ln /etc/init.d/mysql /etc/rc3.d/S91mysql
shell> ln /etc/init.d/mysql /etc/rc0.d/K02mysql
```

To remove MySQL, the installed package name is `mysql`. You can use this in combination with the `pkgrm` command to remove the installation.

To upgrade when using the Solaris package file format, you must remove the existing installation before installing the updated package. Removal of the package does not delete the existing database information, only the server, binaries and support files. The typical upgrade sequence is therefore:

```
shell> mysqladmin shutdown
shell> pkgrm mysql
shell> pkgadd -d mysql-8.0.3-solaris11-x86_64.pkg
shell> mysqld_safe &
shell> mysql_upgrade
```

You should check the notes in Section 2.10, “Upgrading or Downgrading MySQL” before performing any upgrade.

2.7 Installing MySQL on FreeBSD

This section provides information about installing MySQL on variants of FreeBSD Unix.

You can install MySQL on FreeBSD by using the binary distribution provided by Oracle. For more information, see Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

The easiest (and preferred) way to install MySQL is to use the `mysql-server` and `mysql-client` ports available at http://www.freebsd.org/. Using these ports gives you the following benefits:

- A working MySQL with all optimizations enabled that are known to work on your version of FreeBSD.
- Automatic configuration and build.
- Startup scripts installed in `/usr/local/etc/rc.d`.
- The ability to use `pkg_info -L` to see which files are installed.
- The ability to use `pkg_delete` to remove MySQL if you no longer want it on your machine.

The MySQL build process requires GNU make (`gmake`) to work. If GNU `make` is not available, you must install it first before compiling MySQL.

To install using the ports system:

```
# cd /usr/ports/databases/mysql80-server
# make
...
# cd /usr/ports/databases/mysql80-client
# make
...
```

The standard port installation places the server into `/usr/local/libexec/mysqld`, with the startup script for the MySQL server placed in `/usr/local/etc/rc.d/mysql-server`.

Some additional notes on the BSD implementation:

- To remove MySQL after installation using the ports system:

```
# cd /usr/ports/databases/mysql80-server
# make deinstall
...
# cd /usr/ports/databases/mysql80-client
# make deinstall
...
```

- If you get problems with the current date in MySQL, setting the `TZ` variable should help. See Section 4.9, “MySQL Program Environment Variables”.

2.8 Installing MySQL from Source

Building MySQL from the source code enables you to customize build parameters, compiler optimizations, and installation location. For a list of systems on which MySQL is known to run, see http://www.mysql.com/support/supportedplatforms/database.html.

Before you proceed with an installation from source, check whether Oracle produces a precompiled binary distribution for your platform and whether it works for you. We put a great deal of effort into ensuring that our binaries are built with the best possible options for optimal performance. Instructions for installing binary distributions are available in Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

Warning

Building MySQL with nonstandard options may lead to reduced functionality, performance, or security.

The MySQL source code contains internal documentation written using Doxygen. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, “Generating MySQL Doxygen Documentation Content”.

Source Installation Methods

There are two methods for installing MySQL from source:

- Use a standard MySQL source distribution. To obtain a standard distribution, see Section 2.1.2, “How to Get MySQL”. For instructions on building from a standard distribution, see Section 2.8.2, “Installing MySQL Using a Standard Source Distribution”.

  Standard distributions are available as compressed `tar` files, Zip archives, or RPM packages. Distribution files have names of the form `mysql-`*`VERSION`*`.tar.gz`, `mysql-`*`VERSION`*`.zip`, or `mysql-`*`VERSION`*`.rpm`, where *`VERSION`* is a number like `8.0.3`. File names for source distributions can be distinguished from those for precompiled binary distributions in that source distribution names are generic and include no platform name, whereas binary distribution names include a platform name indicating the type of system for which the distribution is intended (for example, `pc-linux-i686` or `winx64`).

- Use a MySQL development tree. For information on building from one of the development trees, see Section 2.8.3, “Installing MySQL Using a Development Source Tree”.

Source Installation System Requirements

Installation of MySQL from source requires several development tools. Some of these tools are needed no matter whether you use a standard source distribution or a development source tree. Other tool requirements depend on which installation method you use.

To install MySQL from source, the following system requirements must be satisfied, regardless of installation method:

- `CMake`, which is used as the build framework on all platforms. `CMake` can be downloaded from http://www.cmake.org.
- A good `make` program. Although some platforms come with their own `make` implementations, it is highly recommended that you use GNU `make` 3.75 or higher. It may already be available on your system as `gmake`. GNU `make` is available from http://www.gnu.org/software/make/.
- MySQL 8.0 source code permits use of C++11 features. To enable a good level of C++11 support across all supported platforms, the following minimum compiler versions apply:
  - GCC: 4.8 or higher
  - Clang: 3.4 or higher (Xcode 7 on macOS)
  - Solaris Studio: 12.4 or higher (Solaris client build only)
  - Visual Studio: 2015
  - CMake: On Windows, the required Visual Studio version results in a required CMake version of 3.2.3 or higher
- The Boost C++ libraries are required to build MySQL (but not to use it). The current version of Boost must be installed. To obtain Boost and its installation instructions, visit the official site. After Boost is installed, tell the build system where the Boost files are located by defining the `WITH_BOOST` option when you invoke `CMake`. For example:

  ```
  shell> cmake . -DWITH_BOOST=/usr/local/boost_1_60_0
  ```

  Adjust the path as necessary to match your installation.
- Sufficient free memory. If you encounter problems such as “internal compiler error” when compiling large source files, it may be that you have too little memory. If compiling on a virtual machine, try increasing the memory allocation.
- Perl is needed if you intend to run test scripts. Most Unix-like systems include Perl. On Windows, you can use a version such as ActiveState Perl.

To install MySQL from a standard source distribution, one of the following tools is required to unpack the distribution file:

- For a `.tar.gz` compressed `tar` file: GNU `gunzip` to uncompress the distribution and a reasonable `tar` to unpack it. If your `tar` program supports the `z` option, it can both uncompress and unpack the file.

  GNU `tar` is known to work. The standard `tar` provided with some operating systems is not able to unpack the long file names in the MySQL distribution. You should download and install GNU `tar`, or if available, use a preinstalled version of GNU tar. Usually this is available as `gnutar`, `gtar`, or as `tar` within a GNU or Free Software directory, such as `/usr/sfw/bin` or `/usr/local/bin`. GNU `tar` is available from http://www.gnu.org/software/tar/.

- For a `.zip` Zip archive: `WinZip` or another tool that can read `.zip` files.
- For an `.rpm` RPM package: The `rpmbuild` program used to build the distribution unpacks it.

To install MySQL from a development source tree, the following additional tools are required:

- The Git revision control system is required to obtain the development source code. The GitHub Help provides instructions for downloading and installing Git on different platforms. MySQL officially joined GitHub in September, 2014. For more information about MySQL's move to GitHub, refer to the announcement on the MySQL Release Engineering blog: MySQL on GitHub
- `bison` 2.1 or higher, available from http://www.gnu.org/software/bison/. (Version 1 is no longer supported.) Use the latest version of `bison` where possible; if you experience problems, upgrade to a later version, rather than revert to an earlier one.

  `bison` is available from http://www.gnu.org/software/bison/. `bison` for Windows can be downloaded from http://gnuwin32.sourceforge.net/packages/bison.htm. Download the package labeled "Complete package, excluding sources". On Windows, the default location for `bison` is the `C:\Program Files\GnuWin32` directory. Some utilities may fail to find `bison` because of the space in the directory name. Also, Visual Studio may simply hang if there are spaces in the path. You can resolve these problems by installing into a directory that does not contain a space; for example `C:\GnuWin32`.
- On Solaris Express, `m4` must be installed in addition to `bison`. `m4` is available from http://www.gnu.org/software/m4/.

Note

If you have to install any programs, modify your `PATH` environment variable to include any directories in which the programs are located. See Section 4.2.10, "Setting Environment Variables".

If you run into problems and need to file a bug report, please use the instructions in Section 1.7, "How to Report Bugs or Problems".

2.8.1 MySQL Layout for Source Installation

By default, when you install MySQL after compiling it from source, the installation step installs files under `/usr/local/mysql`. The component locations under the installation directory are the same as for binary distributions. See Table 2.3, "MySQL Installation Layout for Generic Unix/Linux Binary Package", and Section 2.3.1, "MySQL Installation Layout on Microsoft Windows". To configure installation locations different from the defaults, use the options described at Section 2.8.4, "MySQL Source-Configuration Options".

2.8.2 Installing MySQL Using a Standard Source Distribution

To install MySQL from a standard source distribution:

1. Verify that your system satisfies the tool requirements listed at Section 2.8, "Installing MySQL from Source".
2. Obtain a distribution file using the instructions in Section 2.1.2, "How to Get MySQL".
3. Configure, build, and install the distribution using the instructions in this section.
4. Perform postinstallation procedures using the instructions in Section 2.9, "Postinstallation Setup and Testing".

In MySQL 8.0, `CMake` is used as the build framework on all platforms. The instructions given here should enable you to produce a working installation. For additional information on using `CMake` to build MySQL, see How to Build MySQL Server with CMake.

If you start from a source RPM, use the following command to make a binary RPM that you can install. If you do not have `rpmbuild`, use `rpm` instead.

```
shell> rpmbuild --rebuild --clean MySQL-VERSION.src.rpm
```

The result is one or more binary RPM packages that you install as indicated in Section 2.5.4, “Installing MySQL on Linux Using RPM Packages from Oracle”.

The sequence for installation from a compressed `tar` file or Zip archive source distribution is similar to the process for installing from a generic binary distribution (see Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”), except that it is used on all platforms and includes steps to configure and compile the distribution. For example, with a compressed `tar` file source distribution on Unix, the basic installation command sequence looks like this:

```
# Preconfiguration setup
shell> groupadd mysql
shell> useradd -r -g mysql -s /bin/false mysql
# Beginning of source-build specific instructions
shell> tar zxvf mysql-VERSION.tar.gz
shell> cd mysql-VERSION
shell> mkdir bld
shell> cd bld
shell> cmake ..
shell> make
shell> make install
# End of source-build specific instructions
# Postinstallation setup
shell> cd /usr/local/mysql
shell> chown -R mysql .
shell> chgrp -R mysql .
shell> bin/mysqld --initialize --user=mysql
shell> bin/mysql_ssl_rsa_setup
shell> chown -R root .
shell> chown -R mysql data
shell> bin/mysqld_safe --user=mysql &
# Next command is optional
shell> cp support-files/mysql.server /etc/init.d/mysql.server
```

A more detailed version of the source-build specific instructions is shown following.

Note

The procedure shown here does not set up any passwords for MySQL accounts. After following the procedure, proceed to Section 2.9, “Postinstallation Setup and Testing”, for postinstallation setup and testing.

Perform Preconfiguration Setup

On Unix, set up the `mysql` user and group that will be used to run and execute the MySQL server and own the database directory. For details, see Creating a `mysql` System User and Group, in Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”. Then perform the following steps as the `mysql` user, except as noted.

Obtain and Unpack the Distribution

Pick the directory under which you want to unpack the distribution and change location into it.

Obtain a distribution file using the instructions in Section 2.1.2, “How to Get MySQL”.

Unpack the distribution into the current directory:

- To unpack a compressed `tar` file, `tar` can uncompress and unpack the distribution if it has `z` option support:

```
shell> tar zxvf mysql-VERSION.tar.gz
```

If your `tar` does not have `z` option support, use `gunzip` to unpack the distribution and `tar` to unpack it:

```
shell> gunzip < mysql-VERSION.tar.gz | tar xvf -
```

Alternatively, `CMake` can uncompress and unpack the distribution:

```
shell> cmake -E tar zxvf mysql-VERSION.tar.gz
```

- To unpack a Zip archive, use `WinZip` or another tool that can read `.zip` files.

Unpacking the distribution file creates a directory named `mysql-VERSION`.

Configure the Distribution

Change location into the top-level directory of the unpacked distribution:

```
shell> cd mysql-VERSION
```

Build out of the source tree to keep the tree clean. If the top-level source directory is named `mysql-src` under your current working directory, you can build in a directory named `bld` at the same level. Create the directory and go there:

```
shell> mkdir bld
shell> cd bld
```

Configure the build directory. The minimum configuration command includes no options to override configuration defaults:

```
shell> cmake ../mysql-src
```

The build directory needs not be outside the source tree. For example, you can build in a directory named `bld` under the top-level source tree. To do this, starting with `mysql-src` as your current working directory, create the directory `bld` and then go there:

```
shell> mkdir bld
shell> cd bld
```

Configure the build directory. The minimum configuration command includes no options to override configuration defaults:

```
shell> cmake ..
```

If you have multiple source trees at the same level (for example, to build multiple versions of MySQL), the second strategy can be advantageous. The first strategy places all build directories at the same level, which requires that you choose a unique name for each. With the second strategy, you can use the same name for the build directory within each source tree. The following instructions assume this second strategy.

On Windows, specify the development environment. For example, the following commands configure MySQL for 32-bit or 64-bit builds, respectively:

```
shell> cmake .. -G "Visual Studio 12 2013"
shell> cmake .. -G "Visual Studio 12 2013 Win64"
```

On macOS, to use the Xcode IDE:

```
shell> cmake .. -G Xcode
```

When you run `cmake`, you might want to add options to the command line. Here are some examples:

- `-DBUILD_CONFIG=mysql_release`: Configure the source with the same build options used by Oracle to produce binary distributions for official MySQL releases.
- `-DCMAKE_INSTALL_PREFIX=dir_name`: Configure the distribution for installation under a particular location.
- `-DCPACK_MONOLITHIC_INSTALL=1`: Cause `make package` to generate a single installation file rather than multiple files.
- `-DWITH_DEBUG=1`: Build the distribution with debugging support.

For a more extensive list of options, see Section 2.8.4, “MySQL Source-Configuration Options”.

To list the configuration options, use one of the following commands:

```
shell> cmake .. -L   # overview
shell> cmake .. -LH  # overview with help text
shell> cmake .. -LAH # all params with help text
shell> ccmake ..     # interactive display
```

If `CMake` fails, you might need to reconfigure by running it again with different options. If you do reconfigure, take note of the following:

- If `CMake` is run after it has previously been run, it may use information that was gathered during its previous invocation. This information is stored in `CMakeCache.txt`. When `CMake` starts up, it looks for that file and reads its contents if it exists, on the assumption that the information is still correct. That assumption is invalid when you reconfigure.
- Each time you run `CMake`, you must run `make` again to recompile. However, you may want to remove old object files from previous builds first because they were compiled using different configuration options.

To prevent old object files or configuration information from being used, run these commands in the build direcotry on Unix before re-running `CMake`:

```
shell> make clean
shell> rm CMakeCache.txt
```

Or, on Windows:

```
shell> devenv MySQL.sln /clean
shell> del CMakeCache.txt
```

If you are going to send mail to a MySQL mailing list to ask for configuration assistance, first check the files in the `CMakeFiles` directory for useful information about the failure. To file a bug report, please use the instructions in Section 1.7, “How to Report Bugs or Problems”.

Build the Distribution

On Unix:

```
shell> make
shell> make VERBOSE=1
```

The second command sets `VERBOSE` to show the commands for each compiled source.

Use `gmake` instead on systems where you are using GNU `make` and it has been installed as `gmake`.

On Windows:

```
shell> devenv MySQL.sln /build RelWithDebInfo
```

If you have gotten to the compilation stage, but the distribution does not build, see Section 2.8.5, “Dealing with Problems Compiling MySQL”, for help. If that does not solve the problem, please enter it into our bugs database using the instructions given in Section 1.7, “How to Report Bugs or Problems”. If you have installed the latest versions of the required tools, and they crash trying to process our configuration files, please report that also. However, if you get a `command not found` error or a similar problem for required tools, do not report it. Instead, make sure that all the required tools are installed and that your `PATH` variable is set correctly so that your shell can find them.

Install the Distribution

On Unix:

```
shell> make install
```

This installs the files under the configured installation directory (by default, `/usr/local/mysql`). You might need to run the command as `root`.

To install in a specific directory, add a `DESTDIR` parameter to the command line:

```
shell> make install DESTDIR="/opt/mysql"
```

Alternatively, generate installation package files that you can install where you like:

```
shell> make package
```

This operation produces one or more `.tar.gz` files that can be installed like generic binary distribution packages. See Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”. If you run `CMake` with `-DCPACK_MONOLITHIC_INSTALL=1`, the operation produces a single file. Otherwise, it produces multiple files.

On Windows, generate the data directory, then create a `.zip` archive installation package:

```
shell> devenv MySQL.sln /build RelWithDebInfo /project initial_database
shell> devenv MySQL.sln /build RelWithDebInfo /project package
```

You can install the resulting `.zip` archive where you like. See Section 2.3.5, “Installing MySQL on Microsoft Windows Using a noinstall Zip Archive”.

Perform Postinstallation Setup

The remainder of the installation process involves setting up the configuration file, creating the core databases, and starting the MySQL server. For instructions, see Section 2.9, “Postinstallation Setup and Testing”.

Note

The accounts that are listed in the MySQL grant tables initially have no passwords. After starting the server, you should set up passwords for them using the instructions in Section 2.9, “Postinstallation Setup and Testing”.

2.8.3 Installing MySQL Using a Development Source Tree

This section describes how to install MySQL from the latest development source code, which is hosted on GitHub. To obtain the MySQL Server source code from this repository hosting service, you can set up a local MySQL Git repository.

On GitHub, MySQL Server and other MySQL projects are found on the MySQL page. The MySQL Server project is a single repository that contains branches for several MySQL series.

MySQL officially joined GitHub in September, 2014. For more information about MySQL's move to GitHub, refer to the announcement on the MySQL Release Engineering blog: MySQL on GitHub

Prerequisites for Installing from Development Source

To install MySQL from a development source tree, your system must satisfy the tool requirements outlined in Section 2.8, “Installing MySQL from Source”.

Setting Up a MySQL Git Repository

To set up a MySQL Git repository on your machine, use this procedure:

1. Clone the MySQL Git repository to your machine. The following command clones the MySQL Git repository to a directory named `mysql-server`. The initial download will take some time to complete, depending on the speed of your connection.

```
~$ git clone https://github.com/mysql/mysql-server.git
Cloning into 'mysql-server'...
remote: Counting objects: 1198513, done.
remote: Total 1198513 (delta 0), reused 0 (delta 0), pack-reused 1198513
Receiving objects: 100% (1198513/1198513), 1.01 GiB | 7.44 MiB/s, done.
Resolving deltas: 100% (993200/993200), done.
Checking connectivity... done.
Checking out files: 100% (25510/25510), done.
```

2. When the clone operation completes, the contents of your local MySQL Git repository appear similar to the following:

```
~$ cd mysql-server
~/mysql-server$ ls
BUILD                cmd-line-utils       libservices          sql
CMakeLists.txt       config.h.cmake       man                  sql-common
COPYING              configure.cmake      mysql-test           storage
Docs                 dbug                 mysys                strings
Doxyfile-perfschema  extra                mysys_ssl            support-files
INSTALL              include              packaging            testclients
README               libbinlogevents      plugin               unittest
VERSION              libbinlogstandalone  rapid                vio
client               libevent             regex                win
cmake                libmysql             scripts              zlib
```

3. Use the `git branch -r` command to view the remote tracking branches for the MySQL repository.

```
~/mysql-server$ git branch -r
  origin/5.5
  origin/5.6
  origin/5.7
  origin/8.0
  origin/HEAD -> origin/5.7
  origin/cluster-7.2
  origin/cluster-7.3
  origin/cluster-7.4
  origin/cluster-7.5
```

4. To view the branches that are checked out in your local repository, issue the `git branch` command. When you cloned the MySQL Git repository, the MySQL 5.7 branch was checked out automatically. The asterisk identifies the 5.7 branch as the active branch.

```
~/mysql-server$ git branch
```

```
* 5.7
```

5. To check out a different MySQL branch, run the `git checkout` command, specifying the branch name. For example, to checkout the MySQL 8.0 branch:

```
~/mysql-server$ git checkout 8.0
Checking out files: 100% (9600/9600), done.
Branch 8.0 set up to track remote branch 8.0 from origin.
Switched to a new branch '8.0'
```

6. Run `git branch` to verify that the MySQL 8.0 branch is present. MySQL 8.0, which is the last branch you checked out, is marked by an asterisk indicating that it is the active branch.

```
~/mysql-server$ git branch
  5.7
* 8.0
```

7. Use the `git checkout` command to switch between branches. For example:

```
~/mysql-server$ git checkout 5.7
```

8. To obtain changes made after your initial setup of the MySQL Git repository, switch to the branch you want to update and issue the `git pull` command:

```
~/mysql-server$ git checkout 8.0
~/mysql-server$ git pull
```

   To examine the commit history, use the `git log` option:

```
~/mysql-server$ git log
```

   You can also browse commit history and source code on the GitHub MySQL site.

   If you see changes or code that you have a question about, send an email to the MySQL `internals` mailing list. See Section 1.6.2, “MySQL Mailing Lists”. For information about contributing a patch, see Contributing to MySQL Server.

9. After you have cloned the MySQL Git repository and have checked out the branch you want to build, you can build MySQL Server from the source code. Instructions are provided in Section 2.8.2, “Installing MySQL Using a Standard Source Distribution”, except that you skip the part about obtaining and unpacking the distribution.

   Be careful about installing a build from a distribution source tree on a production machine. The installation command may overwrite your live release installation. If you already have MySQL installed and do not want to overwrite it, run `CMake` with values for the `CMAKE_INSTALL_PREFIX`, `MYSQL_TCP_PORT`, and `MYSQL_UNIX_ADDR` options different from those used by your production server. For additional information about preventing multiple servers from interfering with each other, see Section 5.7, “Running Multiple MySQL Instances on One Machine”.

   Play hard with your new installation. For example, try to make new features crash. Start by running `make test`. See Section 28.1.2, “The MySQL Test Suite”.

2.8.4 MySQL Source-Configuration Options

The `CMake` program provides a great deal of control over how you configure a MySQL source distribution. Typically, you do this using options on the `CMake` command line. For information about options supported by `CMake`, run either of these commands in the top-level source directory:

```
shell> cmake . -LH
shell> ccmake .
```

You can also affect `CMake` using certain environment variables. See Section 4.9, “MySQL Program Environment Variables”.

The following table shows the available `CMake` options. In the `Default` column, `PREFIX` stands for the value of the `CMAKE_INSTALL_PREFIX` option, which specifies the installation base directory. This value is used as the parent location for several of the installation subdirectories.

Table 2.12 MySQL Source-Configuration Option Reference (`CMake`)

Formats	Description	Default	Introduced	Removed
`BUILD_CONFIG`	Use same build options as official releases			
`CMAKE_BUILD_TYPE`	Type of build to produce	`RelWithDebInfo`		
`CMAKE_CXX_FLAGS`	Flags for C++ Compiler			
`CMAKE_C_FLAGS`	Flags for C Compiler			
`CMAKE_INSTALL_PREFIX`	Installation base directory	`/usr/local/mysql`		
`COMPILATION_COMMENT`	Comment about compilation environment			
`CPACK_MONOLITHIC_INSTALL`	Whether package build produces single file	`OFF`		
`DEFAULT_CHARSET`	The default server character set	`latin1`		
`DEFAULT_COLLATION`	The default server collation	`latin1_swedish_ci`		
`DISABLE_PSI_COND`	Exclude Performance Schema condition instrumentation	`OFF`		
`DISABLE_PSI_FILE`	Exclude Performance Schema file instrumentation	`OFF`		
`DISABLE_PSI_IDLE`	Exclude Performance Schema idle instrumentation	`OFF`		
`DISABLE_PSI_MEMORY`	Exclude Performance Schema memory instrumentation	`OFF`		
`DISABLE_PSI_METADATA`	Exclude Performance Schema metadata instrumentation	`OFF`		
`DISABLE_PSI_MUTEX`	Exclude Performance Schema mutex instrumentation	`OFF`		
`DISABLE_PSI_RWLOCK`	Exclude Performance Schema rwlock instrumentation	`OFF`		
`DISABLE_PSI_SOCKET`	Exclude Performance Schema socket instrumentation	`OFF`		
`DISABLE_PSI_SP`	Exclude Performance Schema stored program instrumentation	`OFF`		
`DISABLE_PSI_STAGE`	Exclude Performance Schema stage instrumentation	`OFF`		

Formats	Description	Default	Introduced	Removed
DISABLE_PSI_STATEMENT	Exclude Performance Schema statement instrumentation	OFF		
DISABLE_PSI_STATEMENT_DIGEST	Exclude Performance Schema statements_digest instrumentation	OFF		
DISABLE_PSI_TABLE	Exclude Performance Schema table instrumentation	OFF		
DISABLE_SHARED	Do not build shared libraries, compile position-dependent code	OFF		
-DLINK_RANDOMIZE	Whether to randomize order of symbols in mysqld binary	OFF	8.0.1	
-DLINK_RANDOMIZE_SEED	Seed value for LINK_RANDOMIZE option	mysql	8.0.1	
DOWNLOAD_BOOST	Whether to download the Boost library	OFF		
DOWNLOAD_BOOST_TIMEOUT	Timeout in seconds for downloading the Boost library	600		
-DUSE_LD_GOLD	Whether to use GNU gold loader	ON	8.0.0	
-DWITH_PROTOBUF	Which Protocol Buffers package to use	bundled		
ENABLED_LOCAL_INFILE	Whether to enable LOCAL for LOAD DATA INFILE	OFF		
ENABLED_PROFILING	Whether to enable query profiling code	ON		
ENABLE_DEBUG_SYNC	Whether to enable Debug Sync support	ON		8.0.1
ENABLE_DOWNLOADS	Whether to download optional files	OFF		
ENABLE_DTRACE	Whether to include DTrace support			8.0.1
ENABLE_GCOV	Whether to include gcov support			
ENABLE_GPROF	Enable gprof (optimized Linux builds only)	OFF		
FORCE_UNSUPPORTED_COMPILER	Whether to permit unsupported compiler	OFF		
IGNORE_AIO_CHECK	With -DBUILD_CONFIG=mysql_release, ignore libaio check	OFF		
INSTALL_BINDIR	User executables directory	PREFIX/bin		
INSTALL_DOCDIR	Documentation directory	PREFIX/docs		
INSTALL_DOCREADMEDIR	README file directory	PREFIX		
INSTALL_INCLUDEDIR	Header file directory	PREFIX/include		
INSTALL_INFODIR	Info file directory	PREFIX/docs		

Formats	Description	Default	Introduced	Removed
INSTALL_LAYOUT	Select predefined installation layout	STANDALONE		
INSTALL_LIBDIR	Library file directory	PREFIX/lib		
INSTALL_MANDIR	Manual page directory	PREFIX/man		
INSTALL_MYSQLKEYRINGDIR	Directory for keyring_file plugin data file	platform specific		
INSTALL_MYSQLSHAREDIR	Shared data directory	PREFIX/share		
INSTALL_MYSQLTESTDIR	mysql-test directory	PREFIX/mysql-test		
INSTALL_PKGCONFIGDIR	Directory for mysqlclient.pc pkg-config file	INSTALL_LIBDIR/pkgconfig		
INSTALL_PLUGINDIR	Plugin directory	PREFIX/lib/plugin		
INSTALL_SBINDIR	Server executable directory	PREFIX/bin		
INSTALL_SECURE_FILE_PRIVDIR	secure_file_priv default value	platform specific		
INSTALL_SHAREDIR	aclocal/mysql.m4 installation directory	PREFIX/share		
INSTALL_STATIC_LIBRARIES	Whether to install static libraries	ON		
INSTALL_SUPPORTFILESDIR	Extra support files directory	PREFIX/support-files		
MAX_INDEXES	Maximum indexes per table	64		
MUTEX_TYPE	InnoDB mutex type	event		
MYSQLX_TCP_PORT	TCP/IP port number used by X Plugin	33060		
MYSQLX_UNIX_ADDR	Unix socket file used by X Plugin	/tmp/mysqlx.sock		
MYSQL_DATADIR	Data directory			
MYSQL_MAINTAINER_MODE	Whether to enable MySQL maintainer-specific development environment	OFF		
MYSQL_PROJECT_NAME	Windows/OS X project name	MySQL		
MYSQL_TCP_PORT	TCP/IP port number	3306		
MYSQL_UNIX_ADDR	Unix socket file	/tmp/mysql.sock		
ODBC_INCLUDES	ODBC includes directory			
ODBC_LIB_DIR	ODBC library directory			
OPTIMIZER_TRACE	Whether to support optimizer tracing			
SYSCONFDIR	Option file directory			
SYSTEMD_PID_DIR	Directory for PID file under systemd	/var/run/mysqld		
SYSTEMD_SERVICE_NAME	Name of MySQL service under systemd	mysqld		

Formats	Description	Default	Introduced	Removed
TMPDIR	tmpdir default value			
WIN_DEBUG_NO_INLINE	Whether to disable function inlining	OFF		
WITHOUT_SERVER	Do not build the server	OFF		
WITHOUT_xxx_STORAGE_ENGINE	Exclude storage engine xxx from build			
WITH_ASAN	Enable AddressSanitizer	OFF		
WITH_AUTHENTICATION_LDAP	Whether to report error if LDAP authentication plugins cannot be built	OFF	8.0.2	
WITH_AUTHENTICATION_PAM	Build PAM authentication plugin	OFF		
WITH_BOOST	The location of the Boost library sources			
WITH_CLIENT_PROTOCOL_TRACING	Build client-side protocol tracing framework	ON		
WITH_DEBUG	Whether to include debugging support	OFF		
WITH_DEFAULT_COMPILER_OPTIONS	Whether to use default compiler options	ON		
WITH_DEFAULT_FEATURE_SET	Whether to use default feature set	ON		
WITH_EDITLINE	Which libedit/editline library to use	bundled		
WITH_INNODB_EXTRA_DEBUG	Whether to include extra debugging support for InnoDB.	OFF		
WITH_INNODB_MEMCACHED	Whether to generate memcached shared libraries.	OFF		
WITH_KEYRING_TEST	Build the keyring test program	OFF		
WITH_LIBEVENT	Which libevent library to use	bundled		
WITH_LIBWRAP	Whether to include libwrap (TCP wrappers) support	OFF		
WITH_LZ4	Type of LZ4 support	bundled		
WITH_MECAB	Compiles MeCab			
WITH_MSAN	Enable MemorySanitizer	OFF		
WITH_MSCRT_DEBUG	Enable Visual Studio CRT memory leak tracing	OFF		
WITH_NUMA	Set NUMA memory allocation policy			
WITH_RAPID	Whether to build rapid development cycle plugins	ON		
WITH_SSL	Type of SSL support	bundled		
WITH_SYSTEMD	Enable installation of systemd support files	OFF		

Formats	Description	Default	Introduced	Removed
WITH_TEST_TRACE_PLUGIN	Build test protocol trace plugin	OFF		
WITH_TSAN	Enable ThreadSanitizer	OFF		
WITH_UBSAN	Enable Undefined Behavior Sanitizer	OFF		
WITH_UNIXODBC	Enable unixODBC support	OFF		
WITH_VALGRIND	Whether to compile in Valgrind header files	OFF		
WITH_ZLIB	Type of zlib support	bundled		
WITH_xxx_STORAGE_ENGINE	Compile storage engine xxx statically into server			

The following sections provide more information about CMake options.

- General Options
- Installation Layout Options
- Storage Engine Options
- Feature Options
- Compiler Flags

For boolean options, the value may be specified as 1 or ON to enable the option, or as 0 or OFF to disable the option.

Many options configure compile-time defaults that can be overridden at server startup. For example, the CMAKE_INSTALL_PREFIX, MYSQL_TCP_PORT, and MYSQL_UNIX_ADDR options that configure the default installation base directory location, TCP/IP port number, and Unix socket file can be changed at server startup with the --basedir, --port, and --socket options for mysqld. Where applicable, configuration option descriptions indicate the corresponding mysqld startup option.

General Options

- -DBUILD_CONFIG=mysql_release

  This option configures a source distribution with the same build options used by Oracle to produce binary distributions for official MySQL releases.

- -DCMAKE_BUILD_TYPE=*type*

  The type of build to produce:

  - RelWithDebInfo: Enable optimizations and generate debugging information. This is the default MySQL build type.
  - Debug: Disable optimizations and generate debugging information. This build type is also used if the WITH_DEBUG option is enabled. That is, -DWITH_DEBUG=1 has the same effect as -DCMAKE_BUILD_TYPE=Debug.

- -DCPACK_MONOLITHIC_INSTALL=*bool*

  This option affects whether the make package operation produces multiple installation package files or a single file. If disabled, the operation produces multiple installation package files, which may be useful if you want to install only a subset of a full MySQL installation. If enabled, it produces a single file for installing everything.

Installation Layout Options

The `CMAKE_INSTALL_PREFIX` option indicates the base installation directory. Other options with names of the form `INSTALL_xxx` that indicate component locations are interpreted relative to the prefix and their values are relative pathnames. Their values should not include the prefix.

- `-DCMAKE_INSTALL_PREFIX=dir_name`

  The installation base directory.

  This value can be set at server startup with the `--basedir` option.

- `-DINSTALL_BINDIR=dir_name`

  Where to install user programs.

- `-DINSTALL_DOCDIR=dir_name`

  Where to install documentation.

- `-DINSTALL_DOCREADMEDIR=dir_name`

  Where to install `README` files.

- `-DINSTALL_INCLUDEDIR=dir_name`

  Where to install header files.

- `-DINSTALL_INFODIR=dir_name`

  Where to install Info files.

- `-DINSTALL_LAYOUT=name`

  Select a predefined installation layout:

  - `STANDALONE`: Same layout as used for `.tar.gz` and `.zip` packages. This is the default.
  - `RPM`: Layout similar to RPM packages.
  - `SVR4`: Solaris package layout.
  - `DEB`: DEB package layout (experimental).

  You can select a predefined layout but modify individual component installation locations by specifying other options. For example:

  ```
  shell> cmake . -DINSTALL_LAYOUT=SVR4 -DMYSQL_DATADIR=/var/mysql/data
  ```

  The `INSTALL_LAYOUT` value determines the default value of the `secure_file_priv` system and `keyring_file_data` system variables; see the descriptions of those variables in Section 5.1.5, “Server System Variables”.

- `-DINSTALL_LIBDIR=dir_name`

  Where to install library files.

- `-DINSTALL_MANDIR=dir_name`

  Where to install manual pages.

- `-DINSTALL_MYSQLKEYRINGDIR=dir_path`

The default directory to use as the location of the `keyring_file` plugin data file. The default value is platform specific and depends on the value of the `INSTALL_LAYOUT` `CMake` option; see the description of the `keyring_file_data` system variable in Section 5.1.5, “Server System Variables”.

- `-DINSTALL_MYSQLSHAREDIR=dir_name`

  Where to install shared data files.

- `-DINSTALL_MYSQLTESTDIR=dir_name`

  Where to install the `mysql-test` directory. To suppress installation of this directory, explicitly set the option to the empty value (`-DINSTALL_MYSQLTESTDIR=`).

- `-DINSTALL_PKGCONFIGDIR=dir_name`

  The directory in which to install the `mysqlclient.pc` file for use by `pkg-config`. The default value is `INSTALL_LIBDIR/pkgconfig`, unless `INSTALL_LIBDIR` ends with `/mysql`, in which case that is removed first.

- `-DINSTALL_PLUGINDIR=dir_name`

  The location of the plugin directory.

  This value can be set at server startup with the `--plugin_dir` option.

- `-DINSTALL_SBINDIR=dir_name`

  Where to install the `mysqld` server.

- `-DINSTALL_SECURE_FILE_PRIVDIR=dir_name`

  The default value for the `secure_file_priv` system variable. The default value is platform specific and depends on the value of the `INSTALL_LAYOUT` `CMake` option; see the description of the `secure_file_priv` system variable in Section 5.1.5, “Server System Variables”.

- `-DINSTALL_SHAREDIR=dir_name`

  Where to install `aclocal/mysql.m4`.

- `-DINSTALL_STATIC_LIBRARIES=bool`

  Whether to install static libraries. The default is `ON`. If set to `OFF`, these libraries are not installed: `libmysqlclient.a`, `libmysqlservices.a`.

- `-DINSTALL_SUPPORTFILESDIR=dir_name`

  Where to install extra support files.

- `-DLINK_RANDOMIZE=bool`

  Whether to randomize the order of symbols in the `mysqld` binary. The default is `OFF`. This option should be enabled only for debugging purposes.

- `-DLINK_RANDOMIZE_SEED=val`

  Seed value for the `LINK_RANDOMIZE` option. The value is a string. The default is `mysql`, an arbitrary choice.

- `-DMYSQL_DATADIR=dir_name`

  The location of the MySQL data directory.

This value can be set at server startup with the `--datadir` option.

- `-DODBC_INCLUDES=dir_name`

  The location of the ODBC includes directory, and may be used while configuring Connector/ODBC.

- `-DODBC_LIB_DIR=dir_name`

  The location of the ODBC library directory, and may be used while configuring Connector/ODBC.

- `-DSYSCONFDIR=dir_name`

  The default `my.cnf` option file directory.

  This location cannot be set at server startup, but you can start the server with a given option file using the `--defaults-file=file_name` option, where *`file_name`* is the full path name to the file.

- `-DSYSTEMD_PID_DIR=dir_name`

  The name of the directory in which to create the PID file when MySQL is managed by systemd. The default is `/var/run/mysqld`; this might be changed implicitly according to the `INSTALL_LAYOUT` value.

  This option is ignored unless `WITH_SYSTEMD` is enabled.

- `-DSYSTEMD_SERVICE_NAME=name`

  The name of the MySQL service to use when MySQL is managed by systemd. The default is `mysqld`; this might be changed implicitly according to the `INSTALL_LAYOUT` value.

  This option is ignored unless `WITH_SYSTEMD` is enabled.

- `-DTMPDIR=dir_name`

  The default location to use for the `tmpdir` system variable. If unspecified, the value defaults to `P_tmpdir` in `<stdio.h>`.

Storage Engine Options

Storage engines are built as plugins. You can build a plugin as a static module (compiled into the server) or a dynamic module (built as a dynamic library that must be installed into the server using the `INSTALL PLUGIN` statement or the `--plugin-load` option before it can be used). Some plugins might not support static or dynamic building.

The `InnoDB`, `MyISAM`, `MERGE`, `MEMORY`, and `CSV` engines are mandatory (always compiled into the server) and need not be installed explicitly.

To compile a storage engine statically into the server, use `-DWITH_engine_STORAGE_ENGINE=1`. Some permissible *`engine`* values are `ARCHIVE`, `BLACKHOLE`, `EXAMPLE`, and `FEDERATED`. Examples:

```
-DWITH_ARCHIVE_STORAGE_ENGINE=1
-DWITH_BLACKHOLE_STORAGE_ENGINE=1
```

Note

It is not possible to compile without Performance Schema support. If it is desired to compile without particular types of instrumentation, that can be done with the following `CMake` options:

```
DISABLE_PSI_COND
```

```
DISABLE_PSI_FILE
DISABLE_PSI_IDLE
DISABLE_PSI_MEMORY
DISABLE_PSI_METADATA
DISABLE_PSI_MUTEX
DISABLE_PSI_PS
DISABLE_PSI_RWLOCK
DISABLE_PSI_SOCKET
DISABLE_PSI_SP
DISABLE_PSI_STAGE
DISABLE_PSI_STATEMENT
DISABLE_PSI_STATEMENT_DIGEST
DISABLE_PSI_TABLE
DISABLE_PSI_THREAD
DISABLE_PSI_TRANSACTION
```

For example, to compile without mutex instrumentation, configure MySQL using the `-DDISABLE_PSI_MUTEX=1` option.

To exclude a storage engine from the build, use `-DWITH_engine_STORAGE_ENGINE=0`. Examples:

```
-DWITH_ARCHIVE_STORAGE_ENGINE=0
-DWITH_EXAMPLE_STORAGE_ENGINE=0
-DWITH_FEDERATED_STORAGE_ENGINE=0
```

It is also possible to exclude a storage engine from the build using `-DWITHOUT_engine_STORAGE_ENGINE=1` (but `-DWITH_engine_STORAGE_ENGINE=0` is preferred). Examples:

```
-DWITHOUT_ARCHIVE_STORAGE_ENGINE=1
-DWITHOUT_EXAMPLE_STORAGE_ENGINE=1
-DWITHOUT_FEDERATED_STORAGE_ENGINE=1
```

If neither `-DWITH_engine_STORAGE_ENGINE` nor `-DWITHOUT_engine_STORAGE_ENGINE` are specified for a given storage engine, the engine is built as a shared module, or excluded if it cannot be built as a shared module.

Feature Options

- `-DCOMPILATION_COMMENT=string`

  A descriptive comment about the compilation environment.

- `-DDEFAULT_CHARSET=charset_name`

  The server character set. By default, MySQL uses the `utf8mb4` character set.

  `charset_name` may be one of `binary`, `armscii8`, `ascii`, `big5`, `cp1250`, `cp1251`, `cp1256`, `cp1257`, `cp850`, `cp852`, `cp866`, `cp932`, `dec8`, `eucjpms`, `euckr`, `gb2312`, `gbk`, `geostd8`, `greek`, `hebrew`, `hp8`, `keybcs2`, `koi8r`, `koi8u`, `latin1`, `latin2`, `latin5`, `latin7`, `macce`, `macroman`, `sjis`, `swe7`, `tis620`, `ucs2`, `ujis`, `utf8`, `utf8mb4`, `utf16`, `utf16le`, `utf32`. The permissible character sets are listed in the `cmake/character_sets.cmake` file as the value of `CHARSETS_AVAILABLE`.

  This value can be set at server startup with the `--character_set_server` option.

- `-DDEFAULT_COLLATION=collation_name`

  The server collation. By default, MySQL uses `utf8mb4_0900_ai_ci`. Use the `SHOW COLLATION` statement to determine which collations are available for each character set.

  This value can be set at server startup with the `--collation_server` option.

- `-DDISABLE_PSI_COND=bool`

Whether to exclude the Performance Schema condition instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_FILE=bool`

  Whether to exclude the Performance Schema file instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_IDLE=bool`

  Whether to exclude the Performance Schema idle instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_MEMORY=bool`

  Whether to exclude the Performance Schema memory instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_METADATA=bool`

  Whether to exclude the Performance Schema metadata instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_MUTEX=bool`

  Whether to exclude the Performance Schema mutex instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_RWLOCK=bool`

  Whether to exclude the Performance Schema rwlock instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_SOCKET=bool`

  Whether to exclude the Performance Schema socket instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_SP=bool`

  Whether to exclude the Performance Schema stored program instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_STAGE=bool`

  Whether to exclude the Performance Schema stage instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_STATEMENT=bool`

  Whether to exclude the Performance Schema statement instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_STATEMENT_DIGEST=bool`

  Whether to exclude the Performance Schema statement_digest instrumentation. The default is `OFF` (include).

- `-DDISABLE_PSI_TABLE=bool`

  Whether to exclude the Performance Schema table instrumentation. The default is `OFF` (include).

- `-DDISABLE_SHARED=bool`

  Whether to disable building build shared libraries and compile position-dependent code. The default is `OFF` (compile position-independent code).

  Enabling this option disables `-fPIC`, with the consequence that shared libraries are not built, including plugins and components. Invoking `CMake` produces this warning: `Dynamic plugins are disabled.`

- `-DDOWNLOAD_BOOST=bool`

Whether to download the Boost library. The default is `OFF`.

See the `WITH_BOOST` option for additional discussion about using Boost.

- `-DDOWNLOAD_BOOST_TIMEOUT=seconds`

  The timeout in seconds for downloading the Boost library. The default is 600 seconds.

  See the `WITH_BOOST` option for additional discussion about using Boost.

- `-DENABLE_DEBUG_SYNC=bool`

  As of MySQL 8.0.1, `ENABLE_DEBUG_SYNC` is removed and enabling `WITH_DEBUG` enables Debug Sync.

- `-DENABLE_DOWNLOADS=bool`

  Whether to download optional files. For example, with this option enabled, `CMake` downloads the Google Test distribution that is used by the test suite to run unit tests.

- `-DENABLE_DTRACE=bool`

  Whether to include support for DTrace probes.

  This option was removed in MySQL 8.0.1 because support for DTrace is removed in MySQL 8.1.

- `-DENABLE_GCOV=bool`

  Whether to include gcov support (Linux only).

- `-DENABLE_GPROF=bool`

  Whether to enable `gprof` (optimized Linux builds only).

- `-DENABLED_LOCAL_INFILE=bool`

  This option controls the compiled-in default `LOCAL` capability for the MySQL client library. Clients that make no explicit arrangements therefore have `LOCAL` capability disabled or enabled according to the `ENABLED_LOCAL_INFILE` setting specified at MySQL build time.

  By default, the client library in MySQL binary distributions is compiled with `ENABLED_LOCAL_INFILE` disabled. (Enabled prior to MySQL 8.0.2.) If you compile MySQL from source, configure it with `ENABLED_LOCAL_INFILE` disabled or enabled based on whether clients that make no explicit arrangements should have `LOCAL` capability disabled or enabled, respectively.

  `ENABLED_LOCAL_INFILE` controls the default for client-side `LOCAL` capability. For the server, the `local_infile` system variable controls server-side `LOCAL` capability. To explicitly cause the server to refuse or permit `LOAD DATA LOCAL` statements (regardless of how client programs and libraries are configured at build time or runtime), start `mysqld` with `local_infile` disabled or enabled, respectively. `local_infile` can also be set at runtime. See Section 6.1.6, “Security Issues with LOAD DATA LOCAL”.

- `-DENABLED_PROFILING=bool`

  Whether to enable query profiling code (for the `SHOW PROFILE` and `SHOW PROFILES` statements).

- `-DFORCE_UNSUPPORTED_COMPILER=bool`

  By default, `CMake` checks for minimum versions of supported compilers: Visual Studio 2015 (Windows); GCC 4.8 or Clang 3.4 (Linux); Developer Studio 12.5 (Solaris server); Developer Studio 12.4 or GCC 4.8 (Solaris client library); Clang 3.6 (macOS), Clang 3.4 (FreeBSD). To disable this check, use `-DFORCE_UNSUPPORTED_COMPILER=ON`.

- `-DIGNORE_AIO_CHECK=bool`

  If the `-DBUILD_CONFIG=mysql_release` option is given on Linux, the `libaio` library must be linked in by default. If you do not have `libaio` or do not want to install it, you can suppress the check for it by specifying `-DIGNORE_AIO_CHECK=1`.

- `-DMAX_INDEXES=num`

  The maximum number of indexes per table. The default is 64. The maximum is 255. Values smaller than 64 are ignored and the default of 64 is used.

- `-DMYSQL_MAINTAINER_MODE=bool`

  Whether to enable a MySQL maintainer-specific development environment. If enabled, this option causes compiler warnings to become errors.

- `-DMUTEX_TYPE=type`

  The mutex type used by `InnoDB`. Options include:

  - `event`: Use event mutexes. This is the default value and the original `InnoDB` mutex implementation.
  - `sys`: Use POSIX mutexes on UNIX systems. Use `CRITICAL_SECTION` onjects on Windows, if available.
  - `futex`: Use Linux futexes instead of condition variables to schedule waiting threads.

- `-DMYSQLX_TCP_PORT=port_num`

  The port number on which X Plugin listens for TCP/IP connections. The default is 33060.

  This value can be set at server startup with the `--mysqlx_port` option.

- `-DMYSQLX_UNIX_ADDR=file_name`

  The Unix socket file path on which the server listens for X Plugin socket connections. This must be an absolute path name. The default is `/tmp/mysqlx.sock`.

  This value can be set at server startup with the `--mysqlx-socket` option.

- `-DMYSQL_PROJECT_NAME=name`

  For Windows or macOS, the project name to incorporate into the project file name.

- `-DMYSQL_TCP_PORT=port_num`

  The port number on which the server listens for TCP/IP connections. The default is 3306.

  This value can be set at server startup with the `--port` option.

- `-DMYSQL_UNIX_ADDR=file_name`

  The Unix socket file path on which the server listens for socket connections. This must be an absolute path name. The default is `/tmp/mysql.sock`.

  This value can be set at server startup with the `--socket` option.

- `-DOPTIMIZER_TRACE=bool`

  Whether to support optimizer tracing. See MySQL Internals: Tracing the Optimizer.

- `-DUSE_LD_GOLD=bool`

CMake causes the build process to link with the GNU gold linker if it is available. To suppress use of this linker, specify the -DUSE_LD_GOLD=0 option.

- -DWIN_DEBUG_NO_INLINE=*bool*

  Whether to disable function inlining on Windows. The default is off (inlining enabled).

- -DWITH_ASAN=*bool*

  Whether to enable the AddressSanitizer, for compilers that support it. The default is off.

- -DWITH_AUTHENTICATION_LDAP=*bool*

  Whether to report an error if the LDAP authentication plugins cannot be built:

  - If this option is disabled (the default), the LDAP plugins are built if the required header files and libraries are found. If they are not, CMake displays a note about it.
  - If this option is enabled, a failure to find the required header file andlibraries causes CMake to produce an error, preventing the server from being built.

  This option was added in MySQL 8.0.2.

- -DWITH_AUTHENTICATION_PAM=*bool*

  Whether to build the PAM authentication plugin, for source trees that include this plugin. (See PAM Pluggable Authentication.) If this option is specified and the plugin cannot be compiled, the build fails.

- -DWITH_BOOST=*path_name*

  The Boost library is required to build MySQL. These CMake options enable control over the library source location, and whether to download it automatically:

  - -DWITH_BOOST=*path_name* specifies the Boost library directory location. It is also possible to specify the Boost location by setting the BOOST_ROOT or WITH_BOOST environment variable.

    -DWITH_BOOST=system is also permitted and indicates that the correct version of Boost is installed on the compilation host in the standard location. In this case, the installed version of Boost is used rather than any version included with a MySQL source distribution.

  - -DDOWNLOAD_BOOST=*bool* specifies whether to download the Boost source if it is not present in the specified location. The default is OFF.

  - -DDOWNLOAD_BOOST_TIMEOUT=*seconds* the timeout in seconds for downloading the Boost library. The default is 600 seconds.

  For example, if you normally build MySQL placing the object output in the bld subdirectory of your MySQL source tree, you can build with Boost like this:

```
mkdir bld
cd bld
cmake .. -DDOWNLOAD_BOOST=ON -DWITH_BOOST=$HOME/my_boost
```

  This causes Boost to be downloaded into the my_boost directory under your home directory. If the required Boost version is already there, no download is done. If the required Boost version changes, the newer version is downloaded.

  If Boost is already installed locally and your compiler finds the Boost header files on its own, it may not be necessary to specify the preceding CMake options. However, if the version of Boost required

by MySQL changes and the locally installed version has not been upgraded, you may have build problems. Using the `CMake` options should give you a successful build.

- `-DWITH_CLIENT_PROTOCOL_TRACING=bool`

  Whether to build the client-side protocol tracing framework into the client library. By default, this option is enabled.

  For information about writing protocol trace client plugins, see Section 28.2.4.11, “Writing Protocol Trace Plugins”.

  See also the `WITH_TEST_TRACE_PLUGIN` option.

- `-DWITH_DEBUG=bool`

  Whether to include debugging support.

  Configuring MySQL with debugging support enables you to use the `--debug="d,parser_debug"` option when you start the server. This causes the Bison parser that is used to process SQL statements to dump a parser trace to the server's standard error output. Typically, this output is written to the error log.

  Sync debug checking for the `InnoDB` storage engine is defined under `UNIV_DEBUG` and is available when debugging support is compiled in using the `WITH_DEBUG` option. When debugging support is compiled in, the `innodb_sync_debug` configuration option can be used to enable or disable `InnoDB` sync debug checking.

  Enabling `WITH_DEBUG` also enables Debug Sync. This facility is used for testing and debugging. When compiled in, Debug Sync is disabled by default at runtime. To enable it, start `mysqld` with the `--debug-sync-timeout=N` option, where `N` is a timeout value greater than 0. (The default value is 0, which disables Debug Sync.) `N` becomes the default timeout for individual synchronization points.

  Sync debug checking for the `InnoDB` storage engine is available when debugging support is compiled in using the `WITH_DEBUG` option.

  For a description of the Debug Sync facility and how to use synchronization points, see MySQL Internals: Test Synchronization.

- `-DWITH_DEFAULT_FEATURE_SET=bool`

  Whether to use the flags from `cmake/build_configurations/feature_set.cmake`.

- `-DWITH_EDITLINE=value`

  Which `libedit`/`editline` library to use. The permitted values are `bundled` (the default) and `system`.

- `-DWITH_INNODB_EXTRA_DEBUG=bool`

  Whether to include extra InnoDB debugging support.

  Enabling `WITH_INNODB_EXTRA_DEBUG` turns on extra InnoDB debug checks. This option can only be enabled when `WITH_DEBUG` is enabled.

- `-DWITH_INNODB_MEMCACHED=bool`

  Whether to generate memcached shared libraries (`libmemcached.so` and `innodb_engine.so`).

- `-DWITH_KEYRING_TEST=bool`

  Whether to build the test program that accompanies the `keyring_file` plugin. The default is `OFF`. Test file source code is located in the `plugin/keyring/keyring-test` directory.

- `-DWITH_LIBEVENT=string`

  Which `libevent` library to use. Permitted values are `bundled` (default), `system`, and `yes`. If you specify `system` or `yes`, the system `libevent` library is used if present. If the system library is not found, the bundled `libevent` library is used. The `libevent` library is required by `InnoDB` memcached.

- `-DWITH_LIBWRAP=bool`

  Whether to include `libwrap` (TCP wrappers) support.

- `-DWITH_LZ4=lz4_type`

  The `WITH_LZ4` indicates the source of `zlib` support:

  - `bundled`: Use the `lz4` library bundled with the distribution. This is the default.
  - `system`: Use the system `lz4` library. If `WITH_LZ4` is set to this value, the `lz4_decompress` utility is not built. In this case, the system `lz4` command can be used instead.

- `-DWITH_MSAN=bool`

  Whether to enable MemorySanitizer, for compilers that support it. The default is off.

  For this option to have an effect if enabled, all libraries linked to MySQL must also have been compiled with the option enabled.

- `-DWITH_MECAB={disabled|system|path_name}`

  Use this option to compile the MeCab parser. If you have installed MeCab to its default installation directory, set `-DWITH_MECAB=system`. The `system` option applies to MeCab installations performed from source or from binaries using a native package management utility. If you installed MeCab to a custom installation directory, specify the path to the MeCab installation. For example, `-DWITH_MECAB=/opt/mecab`. If the `system` option does not work, specifying the MeCab installation path should work in all cases.

  For related information, see Section 12.9.9, “MeCab Full-Text Parser Plugin”.

- `-DWITH_MSCRT_DEBUG=bool`

  Whether to enable Visual Studio CRT memory leak tracing. The default is `OFF`.

- `-DWITH_NUMA=bool`

  Explicitly set the NUMA memory allocation policy. `CMake` sets the default `WITH_NUMA` value based on whether the current platform has `NUMA` support. For platforms without NUMA support, `CMake` behaves as follows:

  - With no NUMA option (the normal case), `CMake` continues normally, producing only this warning: NUMA library missing or required version not available
  - With `-DWITH_NUMA=ON`, `CMake` aborts with this error: NUMA library missing or required version not available

- `-DWITH_PROTOBUF=protobuf_type`

  Which Protocol Buffers package to use. *`protobuf_type`* can be one of the following values:

  - `bundled`: Use the package bundled with the distribution. This is the default.
  - `system`: Use the package installed on the system.

  Other values are ignored, with a fallback to `bundled`.

- `-DWITH_RAPID=bool`

  Whether to build the rapid development cycle plugins. When enabled, a `rapid` directory is created in the build tree containing these plugins. When disabled, no `rapid` directory is created in the build tree. The default is `ON`, unless the `rapid` directory is removed from the source tree, in which case the default becomes `OFF`.

- `-DWITH_SSL={ssl_type|path_name}`

  The type of SSL support to include or the path name to the OpenSSL installation to use.

  - *`ssl_type`* can be one of the following values:

    - `yes`: Use the system SSL library if present, else the library bundled with the distribution.
    - `bundled`: Use the SSL library bundled with the distribution. This is the default.
    - `system`: Use the system SSL library.

  - *`path_name`* is the path name to the OpenSSL installation to use. Using this can be preferable to using the *`ssl_type`* value of `system`, for it can prevent CMake from detecting and using an older or incorrect OpenSSL version installed on the system. (Another permitted way to do the same thing is to set the `CMAKE_PREFIX_PATH` option to *`path_name`*.)

  For information about using SSL support, see Section 6.4, “Using Secure Connections”.

- `-DWITH_SYSTEMD=bool`

  Whether to enable installation of systemd support files. By default, this option is disabled. When enabled, systemd support files are installed, and scripts such as `mysqld_safe` and the System V initialization script are not installed. On platforms where systemd is not available, enabling `WITH_SYSTEMD` results in an error from `CMake`.

  For more information about using systemd, see Section 2.5.9, “Managing MySQL Server with systemd”. That section also includes information about specifying options previously specified in `[mysqld_safe]` option groups. Because `mysqld_safe` is not installed when systemd is used, such options must be specified another way.

- `-DWITH_TEST_TRACE_PLUGIN=bool`

  Whether to build the test protocol trace client plugin (see Using the Test Protocol Trace Plugin). By default, this option is disabled. Enabling this option has no effect unless the `WITH_CLIENT_PROTOCOL_TRACING` option is enabled. If MySQL is configured with both options enabled, the `libmysqlclient` client library is built with the test protocol trace plugin built in, and all the standard MySQL clients load the plugin. However, even when the test plugin is enabled, it has no effect by default. Control over the plugin is afforded using environment variables; see Using the Test Protocol Trace Plugin.

  Note

  Do *not* enable the `WITH_TEST_TRACE_PLUGIN` option if you want to use your own protocol trace plugins because only one such plugin can be loaded at a time and an error occurs for attempts to load a second one. If you have already built MySQL with the test protocol trace plugin enabled to see how it works, you must rebuild MySQL without it before you can use your own plugins.

  For information about writing trace plugins, see Section 28.2.4.11, “Writing Protocol Trace Plugins”.

- `-DWITH_TSAN=bool`

Whether to enable the ThreadSanitizer, for compilers that support it. The default is off.

- `-DWITH_UBSAN=bool`

  Whether to enable the Undefined Behavior Sanitizer, for compilers that support it. The default is off.

- `-DWITH_UNIXODBC=1`

  Enables unixODBC support, for Connector/ODBC.

- `-DWITH_VALGRIND=bool`

  Whether to compile in the Valgrind header files, which exposes the Valgrind API to MySQL code. The default is `OFF`.

  To generate a Valgrind-aware debug build, `-DWITH_VALGRIND=1` normally is combined with `-DWITH_DEBUG=1`. See Building Debug Configurations.

- `-DWITH_ZLIB=zlib_type`

  Some features require that the server be built with compression library support, such as the `COMPRESS()` and `UNCOMPRESS()` functions, and compression of the client/server protocol. The `WITH_ZLIB` indicates the source of `zlib` support:

  - `bundled`: Use the `zlib` library bundled with the distribution. This is the default.
  - `system`: Use the system `zlib` library. If `WITH_ZLIB` is set to this value, the `zlib_decompress` utility is not built. In this case, the system `openssl zlib` command can be used instead.

- `-DWITHOUT_SERVER=bool`

  Whether to build without the MySQL server. The default is `OFF`, which does build the server.

Compiler Flags

- `-DCMAKE_C_FLAGS="flags"`

  Flags for the C Compiler.

- `-DCMAKE_CXX_FLAGS="flags"`

  Flags for the C++ Compiler.

- `-DWITH_DEFAULT_COMPILER_OPTIONS=bool`

  Whether to use the flags from `cmake/build_configurations/compiler_options.cmake`.

Note

All optimization flags were carefully chosen and tested by the MySQL build team. Overriding them can lead to unexpected results and is done at your own risk.

To specify your own C and C++ compiler flags, for flags that do not affect optimization, use the `CMAKE_C_FLAGS` and `CMAKE_CXX_FLAGS` CMake options.

When providing your own compiler flags, you might want to specify `CMAKE_BUILD_TYPE` as well.

For example, to create a 32-bit release build on a 64-bit Linux machine, do this:

```
shell> mkdir bld
```

```
shell> cd bld
shell> cmake .. -DCMAKE_C_FLAGS=-m32 \
         -DCMAKE_CXX_FLAGS=-m32 \
         -DCMAKE_BUILD_TYPE=RelWithDebInfo
```

If you set flags that affect optimization (`-O`*number*), you must set the `CMAKE_C_FLAGS_`*build_type* and/or `CMAKE_CXX_FLAGS_`*build_type* options, where *build_type* corresponds to the `CMAKE_BUILD_TYPE` value. To specify a different optimization for the default build type (`RelWithDebInfo`) set the `CMAKE_C_FLAGS_RELWITHDEBINFO` and `CMAKE_CXX_FLAGS_RELWITHDEBINFO` options. For example, to compile on Linux with `-O3` and with debug symbols, do this:

```
shell> cmake .. -DCMAKE_C_FLAGS_RELWITHDEBINFO="-O3 -g" \
         -DCMAKE_CXX_FLAGS_RELWITHDEBINFO="-O3 -g"
```

2.8.5 Dealing with Problems Compiling MySQL

The solution to many problems involves reconfiguring. If you do reconfigure, take note of the following:

- If `CMake` is run after it has previously been run, it may use information that was gathered during its previous invocation. This information is stored in `CMakeCache.txt`. When `CMake` starts up, it looks for that file and reads its contents if it exists, on the assumption that the information is still correct. That assumption is invalid when you reconfigure.
- Each time you run `CMake`, you must run `make` again to recompile. However, you may want to remove old object files from previous builds first because they were compiled using different configuration options.

To prevent old object files or configuration information from being used, run the following commands before re-running `CMake`:

On Unix:

```
shell> make clean
shell> rm CMakeCache.txt
```

On Windows:

```
shell> devenv MySQL.sln /clean
shell> del CMakeCache.txt
```

If you build outside of the source tree, remove and recreate your build directory before re-running `CMake`. For instructions on building outside of the source tree, see How to Build MySQL Server with CMake.

On some systems, warnings may occur due to differences in system include files. The following list describes other problems that have been found to occur most often when compiling MySQL:

- To define which C and C++ compilers to use, you can define the `CC` and `CXX` environment variables. For example:

  ```
  shell> CC=gcc
  shell> CXX=g++
  shell> export CC CXX
  ```

  To specify your own C and C++ compiler flags, use the `CMAKE_C_FLAGS` and `CMAKE_CXX_FLAGS` CMake options. See Compiler Flags.

  To see what flags you might need to specify, invoke `mysql_config` with the `--cflags` and `--cxxflags` options.

- To see what commands are executed during the compile stage, after using `CMake` to configure MySQL, run `make VERBOSE=1` rather than just `make`.
- If compilation fails, check whether the `MYSQL_MAINTAINER_MODE` option is enabled. This mode causes compiler warnings to become errors, so disabling it may enable compilation to proceed.
- If your compile fails with errors such as any of the following, you must upgrade your version of `make` to GNU `make`:

  ```
  make: Fatal error in reader: Makefile, line 18:
  Badly formed macro assignment
  ```

  Or:

  ```
  make: file `Makefile' line 18: Must be a separator (:
  ```

  Or:

  ```
  pthread.h: No such file or directory
  ```

  Solaris and FreeBSD are known to have troublesome `make` programs.

  GNU `make` 3.75 is known to work.

- The `sql_yacc.cc` file is generated from `sql_yacc.yy`. Normally, the build process does not need to create `sql_yacc.cc` because MySQL comes with a pregenerated copy. However, if you do need to re-create it, you might encounter this error:

  ```
  "sql_yacc.yy", line xxx fatal: default action causes potential...
  ```

  This is a sign that your version of `yacc` is deficient. You probably need to install a recent version of `bison` (the GNU version of `yacc`) and use that instead.

  Versions of `bison` older than 1.75 may report this error:

  ```
  sql_yacc.yy:#####: fatal error: maximum table size (32767) exceeded
  ```

  The maximum table size is not actually exceeded; the error is caused by bugs in older versions of `bison`.

For information about acquiring or updating tools, see the system requirements in Section 2.8, “Installing MySQL from Source”.

2.8.6 MySQL Configuration and Third-Party Tools

Third-party tools that need to determine the MySQL version from the MySQL source can read the `VERSION` file in the top-level source directory. The file lists the pieces of the version separately. For example, if the version is MySQL 8.0.4-m14, the file looks like this:

```
MYSQL_VERSION_MAJOR=8
MYSQL_VERSION_MINOR=0
MYSQL_VERSION_PATCH=4
MYSQL_VERSION_EXTRA=-m14
```

If the source is not for a General Availablility (GA) release, the `MYSQL_VERSION_EXTRA` value will be nonempty. For the example, the value corresponds to Milestone 14.

To construct a five-digit number from the version components, use this formula:

```
MYSQL_VERSION_MAJOR*10000 + MYSQL_VERSION_MINOR*100 + MYSQL_VERSION_PATCH
```

2.8.7 Generating MySQL Doxygen Documentation Content

The MySQL source code contains internal documentation written using Doxygen. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the following procedure:

1. Install `doxygen` 1.8.11 or higher. Distributions are available here:

   ```
   http://www.stack.nl/~dimitri/doxygen/download.html
   ```

   After installing `doxygen`, verify the version number:

   ```
   shell> doxygen --version
   1.8.11
   ```

2. Install PlantUML. Distributions are available here:

   ```
   http://plantuml.com/download.html
   ```

   When you install PlantUML on Windows (tested on Windows 10), you must run it at least once as administrator so it creates the registry keys. Open an administrator console and run this command:

   ```
   java -jar path-to-plantuml.jar
   ```

   The command should open a GUI window and return no errors on the console.

3. Install the Graphviz `dot` command. Distributions are available here:

   ```
   http://www.graphviz.org/
   ```

   After installing Graphviz, verify `dot` availability. For example:

   ```
   shell> which dot
   /usr/bin/dot
   shell> dot -V
   dot - graphviz version 2.28.0 (20130928.0220)
   ```

4. Set the `PLANTUML_JAR_PATH` environment to the location where you installed PlantUML. For example:

   ```
   export PLANTUML_JAR_PATH=path-to-plantuml.jar
   ```

5. To invoke `doxygen`, change location to the top-level directory of your MySQL source distribution and execute these commands:

   ```
   mkdir -p generated/doxygen
   doxygen
   ```

   Inspect the error log. It is available in the `doxyerror.log` file in the top-level directory. Assuming that the build executed successfully, view the generated output using a browser. For example:

   ```
   firefox generated/doxygen/html/index.html
   ```

You can use a different browser, and even bookmark the page.

2.9 Postinstallation Setup and Testing

This section discusses tasks that you should perform after installing MySQL:

- If necessary, initialize the data directory and create the MySQL grant tables. For some MySQL installation methods, data directory initialization may be done for you automatically:
  - Windows installation operations performed by MySQL Installer initialize the data directory automatically.
  - Installation on Linux using a server RPM distribution.
  - Installation using the native packaging system on many platforms, including Debian Linux, Ubuntu Linux, Gentoo Linux, and others.
  - Installation on OS X using a DMG distribution.

  For other platforms and installation types, including installation from generic binary and source distributions, you must initialize the data directory yourself. For instructions, see Section 2.9.1, “Initializing the Data Directory”.
- Start the server and make sure that it can be accessed. For instructions, see Section 2.9.2, “Starting the Server”, and Section 2.9.3, “Testing the Server”.
- Assign passwords to the initial `root` account in the grant tables, if that was not already done during data directory initialization. Passwords prevent unauthorized access to the MySQL server. For instructions, see Section 2.9.4, “Securing the Initial MySQL Account”.
- Optionally, arrange for the server to start and stop automatically when your system starts and stops. For instructions, see Section 2.9.5, “Starting and Stopping MySQL Automatically”.
- Optionally, populate time zone tables to enable recognition of named time zones. For instructions, see Section 10.6, “MySQL Server Time Zone Support”.

When you are ready to create additional user accounts, you can find information on the MySQL access control system and account management in Section 6.2, “The MySQL Access Privilege System”, and Section 6.3, “MySQL User Account Management”.

2.9.1 Initializing the Data Directory

After installing MySQL, you must initialize the data directory, including the tables in the `mysql` system database. For some MySQL installation methods, data directory initialization may be done automatically, as described in Section 2.9, “Postinstallation Setup and Testing”. For other installation methods, including installation from generic binary and source distributions, you must initialize the data directory yourself.

This section describes how to initialize the data directory on Unix and Unix-like systems. (For Windows, see Section 2.3.7, “Windows Postinstallation Procedures”.) For some suggested commands that you can use to test whether the server is accessible and working properly, see Section 2.9.3, “Testing the Server”.

In the examples shown here, the server runs under the user ID of the `mysql` login account. This assumes that such an account exists. Either create the account if it does not exist, or substitute the name of a different existing login account that you plan to use for running the server. For information about creating the account, see Creating a `mysql` System User and Group, in Section 2.2, “Installing MySQL on Unix/Linux Using Generic Binaries”.

1. Change location into the top-level directory of your MySQL installation, represented here by *`BASEDIR`*:

```
shell> cd BASEDIR
```

BASEDIR is likely to be something like `/usr/local/mysql` or `/usr/local`. The following steps assume that you have changed location to this directory.

You will find several files and subdirectories in the *BASEDIR* directory. The most important for installation purposes is the `bin` subdirectory, which contains the server as well as client and utility programs.

2. Create a directory that provides a location to use as the value of the `secure_file_priv` system variable that limits import/export operations to a specific directory. See Section 5.1.5, "Server System Variables".

```
shell> mkdir mysql-files
shell> chmod 750 mysql-files
```

3. If necessary, ensure that the distribution contents are accessible to `mysql`. If you installed the distribution as `mysql`, no further action is required. If you installed the distribution as `root`, its contents will be owned by `root`. Change its ownership to `mysql` by executing the following commands as `root` in the installation directory. The first command changes the owner attribute of the files to the `mysql` user. The second changes the group attribute to the `mysql` group.

```
shell> chown -R mysql .
shell> chgrp -R mysql .
```

4. If necessary, initialize the data directory, including the `mysql` database containing the initial MySQL grant tables that determine how users are permitted to connect to the server.

   Typically, data directory initialization need be done only the first time you install MySQL. If you are upgrading an existing installation, you should run `mysql_upgrade` instead (see Section 4.4.5, "`mysql_upgrade` — Check and Upgrade MySQL Tables"). However, the command that initializes the data directory does not overwrite any existing privilege tables, so it should be safe to run in any circumstances.

   Use the server to initialize the data directory:

```
shell> bin/mysqld --initialize --user=mysql
```

   For more information, see Section 2.9.1.1, "Initializing the Data Directory Manually Using mysqld".

5. If you want the server to be able to deploy with automatic support for secure connections, use the `mysql_ssl_rsa_setup` utility to create default SSL and RSA files:

```
shell> mysql_ssl_rsa_setup
```

   For more information, see Section 4.4.3, "`mysql_ssl_rsa_setup` — Create SSL/RSA Files".

6. After initializing the data directory, you can establish the final installation ownership settings. To leave the installation owned by `mysql`, no action is required here. Otherwise, most of the MySQL installation can be owned by `root` if you like. The exception is that the data directory and the `mysql-files` directory must be owned by `mysql`. To accomplish this, run the following commands as `root` in the installation directory. For some distribution types, the data directory might be named `var` rather than `data`; adjust the second command accordingly.

```
shell> chown -R root .
shell> chown -R mysql data mysql-files
```

   If the plugin directory (the directory named by the `plugin_dir` system variable) is writable by the server, it may be possible for a user to write executable code to a file in the directory using `SELECT ... INTO DUMPFILE`. This can be prevented by making the plugin directory read only

to the server or by setting the `secure_file_priv` system variable at server startup to a directory where `SELECT` writes can be performed safely. (For example, set it to the `mysql-files` directory created earlier.)

7. To specify options that the MySQL server should use at startup, put them in a `/etc/my.cnf` or `/etc/mysql/my.cnf` file. You can use such a file, for example, to set the `secure_file_priv` system variable. See Section 5.1.2, “Server Configuration Defaults”. If you do not do this, the server starts with its default settings.

8. If you want MySQL to start automatically when you boot your machine, see Section 2.9.5, “Starting and Stopping MySQL Automatically”.

Data directory initialization creates time zone tables in the `mysql` database but does not populate them. To do so, use the instructions in Section 10.6, “MySQL Server Time Zone Support”.

2.9.1.1 Initializing the Data Directory Manually Using mysqld

This section describes how to initialize the data directory using `mysqld`, the MySQL server.

The following instructions assume that your current location is the MySQL installation directory, represented here by *`BASEDIR`*:

```
shell> cd BASEDIR
```

To initialize the data directory, invoke `mysqld` with the `--initialize` or `--initialize-insecure` option, depending on whether you want the server to generate a random initial password for the `'root'@'localhost'` account.

On Windows, use one of these commands:

```
C:\> bin\mysqld --initialize
C:\> bin\mysqld --initialize-insecure
```

On Unix and Unix-like systems, it is important to make sure that the database directories and files are owned by the `mysql` login account so that the server has read and write access to them when you run it later. To ensure this, start `mysqld` from the system `root` account and include the `--user` option as shown here:

```
shell> bin/mysqld --initialize --user=mysql
shell> bin/mysqld --initialize-insecure --user=mysql
```

Otherwise, execute the program while logged in as `mysql`, in which case you can omit the `--user` option from the command.

Regardless of platform, use `--initialize` for “secure by default” installation (that is, including generation of a random initial `root` password). In this case, the password is marked as expired and you will need to choose a new one. With the `--initialize-insecure` option, no `root` password is generated; it is assumed that you will assign a password to the account in timely fashion before putting the server into production use.

It might be necessary to specify other options such as `--basedir` or `--datadir` if `mysqld` does not identify the correct locations for the installation directory or data directory. For example (enter the command on one line):

```
shell> bin/mysqld --initialize --user=mysql
         --basedir=/opt/mysql/mysql
         --datadir=/opt/mysql/mysql/data
```

Alternatively, put the relevant option settings in an option file and pass the name of that file to `mysqld`. For Unix and Unix-like systems, suppose that the option file name is `/opt/mysql/mysql/etc/my.cnf`. Put these lines in the file:

```
[mysqld]
basedir=/opt/mysql/mysql
datadir=/opt/mysql/mysql/data
```

Then invoke `mysqld` as follows (enter the command on a single line with the `--defaults-file` option first):

```
shell> bin/mysqld --defaults-file=/opt/mysql/mysql/etc/my.cnf
         --initialize --user=mysql
```

On Windows, suppose that `C:\my.ini` contains these lines:

```
[mysqld]
basedir=C:\\Program Files\\MySQL\\MySQL Server 8.0
datadir=D:\\MySQLdata
```

Then invoke `mysqld` as follows (the `--defaults-file` option must be first):

```
C:\> bin/mysqld --defaults-file=C:\my.ini --initialize
```

When invoked with the `--initialize` or `--initialize-insecure` option, `mysqld` performs the following initialization sequence.

Note

The server writes any messages to its standard error output. This may be redirected to the error log, so look there if you do not see the messages on your screen. For information about the error log, including where it is located, see Section 5.4.2, “The Error Log”.

On Windows, use the `--console` option to direct messages to the console.

1. The server checks the existence of the data directory:

   - If no data directory exists, the server creates it.
   - If a data directory exists but is not empty (that is, it contains files or subdirectories), the server exits after producing an error message:

     ```
     [ERROR] --initialize specified but the data directory exists. Aborting.
     ```

     In this case, remove or rename the data directory and try again.

     An existing data directory is permitted to be nonempty if every entry has a name that begins with a period (`.`).

2. Within the data directory, the server creates the `mysql` system database and its tables, including the grant tables, server-side help tables, and time zone tables. For a complete listing and description of the grant tables, see Section 6.2, “The MySQL Access Privilege System”.

3. The server initializes the system tablespace and related data structures needed to manage `InnoDB` tables.

Note

After `mysqld` sets up the `InnoDB` system tablespace, changes to some tablespace characteristics require setting up a whole new instance. This includes the file name of the first file in the system tablespace and the number of undo logs. If you do not want to use the default values, make sure that the settings for the `innodb_data_file_path` and `innodb_log_file_size` configuration parameters are in place

> in the MySQL configuration file before running `mysqld`. Also make sure to specify as necessary other parameters that affect the creation and location of `InnoDB` files, such as `innodb_data_home_dir` and `innodb_log_group_home_dir`.
>
> If those options are in your configuration file but that file is not in a location that MySQL reads by default, specify the file location using the `--defaults-extra-file` option when you run `mysqld`.

4. The server creates a `'root'@'localhost'` superuser account and other reserved accounts (see Section 6.3.5, “Reserved User Accounts”). Some reserved accounts are locked and cannot be used by clients, but `'root'@'localhost'` is intended for administrative use and you should assign it a password.

   The server's action with respect to a password for the `'root'@'localhost'` account depends on how you invoke it:

   - With `--initialize` but not `--initialize-insecure`, the server generates a random password, marks it as expired, and writes a message displaying the password:

     ```
     [Warning] A temporary password is generated for root@localhost:
     iTag*AfrH5ej
     ```

   - With `--initialize-insecure`, (either with or without `--initialize` because `--initialize-insecure` implies `--initialize`), the server does not generate a password or mark it expired, and writes a warning message:

     ```
     Warning] root@localhost is created with an empty password ! Please
     consider switching off the --initialize-insecure option.
     ```

   See later in this section for instructions on assigning a new `'root'@'localhost'` password.

5. The server populates the server-side help tables if content is available (in the `fill_help_tables.sql` file). The server does not populate the time zone tables; to do so, see Section 10.6, “MySQL Server Time Zone Support”.

6. If the `--init-file` option was given to name a file of SQL statements, the server executes the statements in the file. This option enables you to perform custom bootstrapping sequences.

   When the server operates in bootstrap mode, some functionality is unavailable that limits the statements permitted in the file. These include statements that relate to account management (such as `CREATE USER` or `GRANT`), replication, and global transaction identifiers.

7. The server exits.

After you initialize the data directory by starting the server with `--initialize` or `--initialize-insecure`, start the server normally (that is, without either of those options) and assign the `'root'@'localhost'` account a new password:

1. Start the server. For instructions, see Section 2.9.2, “Starting the Server”.

2. Connect to the server:

   - If you used `--initialize` but not `--initialize-insecure` to initialize the data directory, connect to the server as `root` using the random password that the server generated during the initialization sequence:

     ```
     shell> mysql -u root -p
     Enter password: (enter the random root password here)
     ```

     Look in the server error log if you do not know this password.

- If you used `--initialize-insecure` to initialize the data directory, connect to the server as `root` without a password:

```
shell> mysql -u root --skip-password
```

3. After connecting, assign a new `root` password:

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'new_password';
```

Note

The data directory initialization sequence performed by the server does not substitute for the actions performed by `mysql_secure_installation` or `mysql_ssl_rsa_setup`. See Section 4.4.2, “`mysql_secure_installation` — Improve MySQL Installation Security”, and Section 4.4.3, “`mysql_ssl_rsa_setup` — Create SSL/RSA Files”.

2.9.2 Starting the Server

This section describes how start the server on Unix and Unix-like systems. (For Windows, see Section 2.3.5.5, “Starting the Server for the First Time”.) For some suggested commands that you can use to test whether the server is accessible and working properly, see Section 2.9.3, “Testing the Server”.

Start the MySQL server like this if your installation includes `mysqld_safe`:

```
shell> bin/mysqld_safe --user=mysql &
```

Note

For Linux systems on which MySQL is installed using RPM packages, server startup and shutdown is managed using systemd rather than `mysqld_safe`, and `mysqld_safe` is not installed.

Start the server like this if your installation includes systemd support:

```
shell> systemctl start mysqld
```

Substitute the appropriate service name if it differs from `mysqld`; for example, `mysql` on SLES systems.

It is important that the MySQL server be run using an unprivileged (non-`root`) login account. To ensure this, run `mysqld_safe` as `root` and include the `--user` option as shown. Otherwise, you should execute the program while logged in as `mysql`, in which case you can omit the `--user` option from the command.

For further instructions for running MySQL as an unprivileged user, see Section 6.1.5, “How to Run MySQL as a Normal User”.

If the command fails immediately and prints `mysqld ended`, look for information in the error log (which by default is the *`host_name`*`.err` file in the data directory).

If the server is unable to access the data directory it starts or read the grant tables in the `mysql` database, it writes a message to its error log. Such problems can occur if you neglected to create the grant tables by initializing the data directory before proceeding to this step, or if you ran the command that initializes the data directory without the `--user` option. Remove the `data` directory and run the command with the `--user` option.

If you have other problems starting the server, see Section 2.9.2.1, “Troubleshooting Problems Starting the MySQL Server”. For more information about `mysqld_safe`, see Section 4.3.2, “`mysqld_safe` — MySQL Server Startup Script”. For more information about systemd support, see Section 2.5.9, “Managing MySQL Server with systemd”.

2.9.2.1 Troubleshooting Problems Starting the MySQL Server

This section provides troubleshooting suggestions for problems starting the server. For additional suggestions for Windows systems, see Section 2.3.6, “Troubleshooting a Microsoft Windows MySQL Server Installation”.

If you have problems starting the server, here are some things to try:

- Check the error log to see why the server does not start. Log files are located in the data directory (typically `C:\Program Files\MySQL\MySQL Server 8.0\data` on Windows, `/usr/local/mysql/data` for a Unix/Linux binary distribution, and `/usr/local/var` for a Unix/Linux source distribution). Look in the data directory for files with names of the form *`host_name`*`.err` and *`host_name`*`.log`, where *`host_name`* is the name of your server host. Then examine the last few lines of these files. Use `tail` to display them:

  ```
  shell> tail host_name.err
  shell> tail host_name.log
  ```

- Specify any special options needed by the storage engines you are using. You can create a `my.cnf` file and specify startup options for the engines that you plan to use. If you are going to use storage engines that support transactional tables (`InnoDB`, `NDB`), be sure that you have them configured the way you want before starting the server. If you are using `InnoDB` tables, see Section 15.6, “InnoDB Configuration” for guidelines and Section 15.13, “InnoDB Startup Options and System Variables” for option syntax.

  Although storage engines use default values for options that you omit, Oracle recommends that you review the available options and specify explicit values for any options whose defaults are not appropriate for your installation.

- Make sure that the server knows where to find the data directory. The `mysqld` server uses this directory as its current directory. This is where it expects to find databases and where it expects to write log files. The server also writes the pid (process ID) file in the data directory.

  The default data directory location is hardcoded when the server is compiled. To determine what the default path settings are, invoke `mysqld` with the `--verbose` and `--help` options. If the data directory is located somewhere else on your system, specify that location with the `--datadir` option to `mysqld` or `mysqld_safe`, on the command line or in an option file. Otherwise, the server will not work properly. As an alternative to the `--datadir` option, you can specify `mysqld` the location of the base directory under which MySQL is installed with the `--basedir`, and `mysqld` looks for the `data` directory there.

  To check the effect of specifying path options, invoke `mysqld` with those options followed by the `--verbose` and `--help` options. For example, if you change location into the directory where `mysqld` is installed and then run the following command, it shows the effect of starting the server with a base directory of `/usr/local`:

  ```
  shell> ./mysqld --basedir=/usr/local --verbose --help
  ```

  You can specify other options such as `--datadir` as well, but `--verbose` and `--help` must be the last options.

  Once you determine the path settings you want, start the server without `--verbose` and `--help`.

  If `mysqld` is currently running, you can find out what path settings it is using by executing this command:

```
shell> mysqladmin variables
```

Or:

```
shell> mysqladmin -h host_name variables
```

host_name is the name of the MySQL server host.

- Make sure that the server can access the data directory. The ownership and permissions of the data directory and its contents must allow the server to read and modify them.

  If you get `Errcode 13` (which means `Permission denied`) when starting `mysqld`, this means that the privileges of the data directory or its contents do not permit server access. In this case, you change the permissions for the involved files and directories so that the server has the right to use them. You can also start the server as `root`, but this raises security issues and should be avoided.

  Change location into the data directory and check the ownership of the data directory and its contents to make sure the server has access. For example, if the data directory is `/usr/local/mysql/var`, use this command:

  ```
  shell> ls -la /usr/local/mysql/var
  ```

  If the data directory or its files or subdirectories are not owned by the login account that you use for running the server, change their ownership to that account. If the account is named `mysql`, use these commands:

  ```
  shell> chown -R mysql /usr/local/mysql/var
  shell> chgrp -R mysql /usr/local/mysql/var
  ```

  Even with correct ownership, MySQL might fail to start up if there is other security software running on your system that manages application access to various parts of the file system. In this case, reconfigure that software to enable `mysqld` to access the directories it uses during normal operation.

- Verify that the network interfaces the server wants to use are available.

  If either of the following errors occur, it means that some other program (perhaps another `mysqld` server) is using the TCP/IP port or Unix socket file that `mysqld` is trying to use:

  ```
  Can't start server: Bind on TCP/IP port: Address already in use
  Can't start server: Bind on unix socket...
  ```

  Use `ps` to determine whether you have another `mysqld` server running. If so, shut down the server before starting `mysqld` again. (If another server is running, and you really want to run multiple servers, you can find information about how to do so in Section 5.7, “Running Multiple MySQL Instances on One Machine”.)

  If no other server is running, execute the command `telnet your_host_name tcp_ip_port_number`. (The default MySQL port number is 3306.) Then press Enter a couple of times. If you do not get an error message like `telnet: Unable to connect to remote host: Connection refused`, some other program is using the TCP/IP port that `mysqld` is trying to use. Track down what program this is and disable it, or tell `mysqld` to listen to a different port with the `--port` option. In this case, specify the same non-default port number for client programs when connecting to the server using TCP/IP.

  Another reason the port might be inaccessible is that you have a firewall running that blocks connections to it. If so, modify the firewall settings to permit access to the port.

If the server starts but you cannot connect to it, make sure that you have an entry in `/etc/hosts` that looks like this:

```
127.0.0.1       localhost
```

- If you cannot get `mysqld` to start, try to make a trace file to find the problem by using the `--debug` option. See Section 28.5.3, “The DBUG Package”.

2.9.3 Testing the Server

After the data directory is initialized and you have started the server, perform some simple tests to make sure that it works satisfactorily. This section assumes that your current location is the MySQL installation directory and that it has a `bin` subdirectory containing the MySQL programs used here. If that is not true, adjust the command path names accordingly.

Alternatively, add the `bin` directory to your `PATH` environment variable setting. That enables your shell (command interpreter) to find MySQL programs properly, so that you can run a program by typing only its name, not its path name. See Section 4.2.10, “Setting Environment Variables”.

Use `mysqladmin` to verify that the server is running. The following commands provide simple tests to check whether the server is up and responding to connections:

```
shell> bin/mysqladmin version
shell> bin/mysqladmin variables
```

If you cannot connect to the server, specify a `-u root` option to connect as `root`. If you have assigned a password for the `root` account already, you'll also need to specify `-p` on the command line and enter the password when prompted. For example:

```
shell> bin/mysqladmin -u root -p version
Enter password: (enter root password here)
```

The output from `mysqladmin version` varies slightly depending on your platform and version of MySQL, but should be similar to that shown here:

```
shell> bin/mysqladmin version
mysqladmin  Ver 14.12 Distrib 8.0.3, for pc-linux-gnu on i686
...

Server version          8.0.3
Protocol version        10
Connection              Localhost via UNIX socket
UNIX socket             /var/lib/mysql/mysql.sock
Uptime:                 14 days 5 hours 5 min 21 sec

Threads: 1  Questions: 366  Slow queries: 0
Opens: 0  Flush tables: 1  Open tables: 19
Queries per second avg: 0.000
```

To see what else you can do with `mysqladmin`, invoke it with the `--help` option.

Verify that you can shut down the server (include a `-p` option if the `root` account has a password already):

```
shell> bin/mysqladmin -u root shutdown
```

Verify that you can start the server again. Do this by using `mysqld_safe` or by invoking `mysqld` directly. For example:

```
shell> bin/mysqld_safe --user=mysql &
```

If `mysqld_safe` fails, see Section 2.9.2.1, “Troubleshooting Problems Starting the MySQL Server”.

Run some simple tests to verify that you can retrieve information from the server. The output should be similar to that shown here.

Use `mysqlshow` to see what databases exist:

```
shell> bin/mysqlshow
+--------------------+
|     Databases      |
+--------------------+
| information_schema |
| mysql              |
| performance_schema |
| sys                |
+--------------------+
```

The list of installed databases may vary, but will always include the minimum of `mysql` and `information_schema`.

If you specify a database name, `mysqlshow` displays a list of the tables within the database:

```
shell> bin/mysqlshow mysql
Database: mysql
+---------------------------+
|          Tables           |
+---------------------------+
| columns_priv              |
| db                        |
| engine_cost               |
| event                     |
| func                      |
| general_log               |
| gtid_executed             |
| help_category             |
| help_keyword              |
| help_relation             |
| help_topic                |
| innodb_index_stats        |
| innodb_table_stats        |
| ndb_binlog_index          |
| plugin                    |
| proc                      |
| procs_priv                |
| proxies_priv              |
| server_cost               |
| servers                   |
| slave_master_info         |
| slave_relay_log_info      |
| slave_worker_info         |
| slow_log                  |
| tables_priv               |
| time_zone                 |
| time_zone_leap_second     |
| time_zone_name            |
| time_zone_transition      |
| time_zone_transition_type |
| user                      |
+---------------------------+
```

Use the `mysql` program to select information from a table in the `mysql` database:

```
shell> bin/mysql -e "SELECT User, Host, plugin FROM mysql.user" mysql
+------+-----------+-----------------------+
| User | Host      | plugin                |
+------+-----------+-----------------------+
| root | localhost | mysql_native_password |
+------+-----------+-----------------------+
```

At this point, your server is running and you can access it. To tighten security if you have not yet assigned a password to the initial account, follow the instructions in Section 2.9.4, “Securing the Initial MySQL Account”.

For more information about `mysql`, `mysqladmin`, and `mysqlshow`, see Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”, Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”, and Section 4.5.7, “`mysqlshow` — Display Database, Table, and Column Information”.

2.9.4 Securing the Initial MySQL Account

The MySQL installation process involves initializing the data directory, including the `mysql` database containing the grant tables that define MySQL accounts. For details, see Section 2.9, “Postinstallation Setup and Testing”.

This section describes how to assign a password to the initial `root` account created during the MySQL installation procedure, if you have not already done so.

Note

On Windows, you can also perform the process described in this section during installation with MySQL Installer (see Section 2.3.3, “MySQL Installer for Windows”). On all platforms, the MySQL distribution includes `mysql_secure_installation`, a command-line utility that automates much of the process of securing a MySQL installation. MySQL Workbench is available on all platforms, and also offers the ability to manage user accounts (see Chapter 30, *MySQL Workbench*).

A password may already be assigned to the initial account under these circumstances:

- On Windows, installations performed using MySQL Installer give you the option of assigning a password.
- Installation using the OS X installer generates an initial random password, which the installer displays to the user in a dialog box.
- Installation using RPM packages generates an initial random password, which is written to the server error log.
- For data directory initialization performed manually using `mysqld --initialize`, `mysqld` generates an initial random password, marks it expired, and writes it to the server error log. For details, see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

The `mysql.user` grant table defines the initial MySQL user account and its access privileges. Installation of MySQL creates only a `'root'@'localhost'` superuser account that has all privileges and can do anything. If the `root` account has an empty password, your MySQL installation is unprotected: Anyone can connect to the MySQL server as `root` *without a password* and be granted all privileges.

The `'root'@'localhost'` account also has a row in the `mysql.proxies_priv` table that enables granting the `PROXY` privilege for `''@''`, that is, for all users and all hosts. This enables `root` to set up proxy users, as well as to delegate to other accounts the authority to set up proxy users. See Section 6.3.11, “Proxy Users”.

To assign a password for the initial MySQL `root` account, use the following procedure. Replace *`new_password`* in the examples with the password that you want to use.

Start the server if it is not running. For instructions, see Section 2.9.2, “Starting the Server”.

The initial `root` account may or may not have a password. Choose whichever of the following procedures applies:

- If the `root` account exists with an initial random password that has been expired, connect to the server as `root` using that password, then choose a new password. This is the case if the data

directory was initialized using `mysqld --initialize`, either manually or using an installer that does not give you the option of specifying a password during the install operation. Because the password exists, you must use it to connect to the server. But because the password is expired, you cannot use the account for any purpose other than to choose a new password, until you do choose one.

1. If you do not know the initial random password, look in the server error log.
2. Connect to the server as `root` using the password:

```
shell> mysql -u root -p
Enter password: (enter the random root password here)
```

3. Choose a new password to replace the random password:

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'new_password';
```

- If the `root` account exists but has no password, connect to the server as `root` using no password, then assign a password. This is the case if you initialized the data directory using `mysqld --initialize-insecure`.

1. Connect to the server as `root` using no password:

```
shell> mysql -u root --skip-password
```

2. Assign a password:

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'new_password';
```

After assigning the `root` account a password, you must supply that password whenever you connect to the server using the account. For example, to connect to the server using the `mysql` client, use this command:

```
shell> mysql -u root -p
Enter password: (enter root password here)
```

To shut down the server with `mysqladmin`, use this command:

```
shell> mysqladmin -u root -p shutdown
Enter password: (enter root password here)
```

Note

For additional information about setting passwords, see Section 6.3.7, “Assigning Account Passwords”. If you forget your `root` password after setting it, see Section B.5.3.2, “How to Reset the Root Password”.

To set up additional accounts, see Section 6.3.2, “Adding User Accounts”.

2.9.5 Starting and Stopping MySQL Automatically

This section discusses methods for starting and stopping the MySQL server.

Generally, you start the `mysqld` server in one of these ways:

- Invoke `mysqld` directly. This works on any platform.
- On Windows, you can set up a MySQL service that runs automatically when Windows starts. See Section 2.3.5.8, “Starting MySQL as a Windows Service”.

- On Unix and Unix-like systems, you can invoke `mysqld_safe`, which tries to determine the proper options for `mysqld` and then runs it with those options. See Section 4.3.2, “`mysqld_safe` — MySQL Server Startup Script”.
- On Linux systems that support systemd, you can use it to control the server. See Section 2.5.9, “Managing MySQL Server with systemd”.
- On systems that use System V-style run directories (that is, `/etc/init.d` and run-level specific directories), invoke `mysql.server`. This script is used primarily at system startup and shutdown. It usually is installed under the name `mysql`. The `mysql.server` script starts the server by invoking `mysqld_safe`. See Section 4.3.3, “`mysql.server` — MySQL Server Startup Script”.
- On OS X, install a launchd daemon to enable automatic MySQL startup at system startup. The daemon starts the server by invoking `mysqld_safe`. For details, see Section 2.4.3, “Installing a MySQL Launch Daemon”. A MySQL Preference Pane also provides control for starting and stopping MySQL through the System Preferences. See Section 2.4.4, “Installing and Using the MySQL Preference Pane”.
- On Solaris, use the service management framework (SMF) system to initiate and control MySQL startup.

systemd, the `mysqld_safe` and `mysql.server` scripts, Solaris SMF, and the OS X Startup Item (or MySQL Preference Pane) can be used to start the server manually, or automatically at system startup time. systemd, `mysql.server`, and the Startup Item also can be used to stop the server.

The following table shows which option groups the server and startup scripts read from option files.

Table 2.13 MySQL Startup Scripts and Supported Server Option Groups

Script	Option Groups
`mysqld`	`[mysqld]`, `[server]`, `[mysqld-`*`major_version`*`]`
`mysqld_safe`	`[mysqld]`, `[server]`, `[mysqld_safe]`
`mysql.server`	`[mysqld]`, `[mysql.server]`, `[server]`

`[mysqld-`*`major_version`*`]` means that groups with names like `[mysqld-5.7]` and `[mysqld-8.0]` are read by servers having versions 5.7.x, 8.0.x, and so forth. This feature can be used to specify options that can be read only by servers within a given release series.

For backward compatibility, `mysql.server` also reads the `[mysql_server]` group and `mysqld_safe` also reads the `[safe_mysqld]` group. To be current, you should update your option files to use the `[mysql.server]` and `[mysqld_safe]` groups instead.

For more information on MySQL configuration files and their structure and contents, see Section 4.2.6, “Using Option Files”.

2.10 Upgrading or Downgrading MySQL

This section describes the steps to upgrade or downgrade a MySQL installation.

Upgrading is a common procedure, as you pick up bug fixes within the same MySQL release series or significant features between major MySQL releases. You perform this procedure first on some test systems to make sure everything works smoothly, and then on the production systems.

Downgrading is less common. Typically, you undo an upgrade because of some compatibility or performance issue that occurs on a production system, and was not uncovered during initial upgrade verification on the test systems. As with the upgrade procedure, perform and verify the downgrade procedure on some test systems first, before using it on a production system.

2.10.1 Upgrading MySQL

MySQL Server 8.0 incorporates a global data dictionary containing information about database objects in transactional tables. In previous MySQL series, dictionary data was stored in metadata files and nontransactional system tables. This section describes how to upgrade your data directory from the file-based structure to the data-dictionary structure.

Important

The process described here requires that you upgrade to MySQL 8.0 from a MySQL 5.7 GA version (5.7.9 or higher). Upgrades from non-GA versions of MySQL 5.7 or from MySQL series preceding 5.7 are not supported.

- Supported Upgrade Methods
- Supported Upgrade Paths
- Before You Begin
- Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation
- Performing an In-Place Upgrade
- Performing a Logical Upgrade
- Upgrade Troubleshooting

After you upgrade, a data dictionary-enabled server entails some general operational differences; see Section 14.6, “Data Dictionary Usage Differences”.

Note

In the following discussion, MySQL commands that must be run using a MySQL account with administrative privileges include `-u root` on the command line to specify the MySQL `root` user. Commands that require a password for `root` also include a `-p` option. Because `-p` is followed by no option value, such commands prompt for the password. Type the password when prompted and press Enter.

SQL statements can be executed using the `mysql` command-line client (connect as `root` to ensure that you have the necessary privileges).

Supported Upgrade Methods

Supported upgrade methods include:

- *In-Place Upgrade*: Involves shutting down the old MySQL version, replacing the old MySQL binaries or packages with the new ones, restarting MySQL on the existing data directory, and running `mysql_upgrade`.
- *Logical Upgrade*: Involves exporting existing data from the old MySQL version using `mysqldump`, installing the new MySQL version, loading the dump file into the new MySQL version, and running `mysql_upgrade`.

For in-place and logical upgrade procedures, see Performing an In-Place Upgrade, and Performing a Logical Upgrade.

If you run MySQL Server on Windows, refer to the upgrade procedure described in Section 2.3.8, “Upgrading MySQL on Windows”.

If your current MySQL installation was installed on an Enterprise Linux platform or Fedora using the MySQL Yum Repository, see Section 2.10.1.2, “Upgrading MySQL with the MySQL Yum Repository”.

If your current MySQL installation was installed on Ubuntu using the MySQL APT repository, see Section 2.10.1.3, “Upgrading MySQL with the MySQL APT Repository”.

Supported Upgrade Paths

Unless otherwise documented, the following upgrade paths are supported:

- Upgrading from a GA release series version to a newer release series version is supported. Skipping release series versions is also supported. The are currently no GA status releases available for MySQL 8.0.
- Upgrading one release level is supported. For example, upgrading from 5.7 to 8.0 is supported. Upgrading to the latest release series version is recommended before upgrading to the next release level. For example, upgrade to the latest 5.7 release before upgrading to 8.0.

Important

For MySQL 8.0, it is required that you upgrade from a MySQL 5.7 GA version (5.7.9 or higher). Upgrades from non-GA versions of MySQL 5.7 are not supported.

- Upgrading more than one release level is supported, but only if you upgrade one release level at a time. For example, if you currently are running MySQL 5.6 and wish to upgrade to a newer series, upgrade to MySQL 5.7 first before upgrading to MySQL 8.0, and so forth. For information on upgrading to MySQL 5.7 see the *MySQL 5.7 Reference Manual*.
- Direct upgrades that skip a release level (for example, upgrading directly from MySQL 5.6 to 8.0) are not recommended or supported.

The following conditions apply to all upgrade paths:

- Upgrades between General Availability (GA) status releases are supported.
- Upgrades between milestone releases (or from a milestone release to a GA release) are not supported. For example, upgrading from 8.0.0 to 8.0.1 is not supported, as neither are GA status releases.
- For upgrades between versions of a MySQL release series that has reached GA status, you can move the MySQL data files between different versions on systems with the same architecture. This is not necessarily true for upgrades between milestone releases. Use of milestone releases is at your own risk.

Before You Begin

Before upgrading, review the following information and perform the recommended steps:

- Before upgrading, protect your data by creating a backup of your current databases and log files. The backup should include the `mysql` system database, which contains the MySQL data dictionary tables and system tables. See Section 7.2, “Database Backup Methods”.
- Review the Release Notes which provide information about features that are new in the MySQL 8.0 or differ from those found in earlier MySQL releases. Some of these changes may result in incompatibilities.

  For a description of MySQL server features that have been removed in MySQL 8.0, see Features Removed in MySQL 8.0. An upgrade requires changes with respect to those features if you use any of them.

  For listings of MySQL server variables and options that have been added, deprecated, or removed in MySQL 8.0, see Section 1.5, “Server and Status Variables and Options Added, Deprecated, or Removed in MySQL 8.0”. If you use any of these items, an upgrade requires configuration changes.
- Review Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0”. This section describes changes that may require action before or after upgrading.
- If you use replication, review Section 18.4.3, “Upgrading a Replication Setup”.

- If you use XA transactions with `InnoDB`, run `XA RECOVER` before upgrading to check for uncommitted XA transactions. If results are returned, either commit or rollback the XA transactions by issuing an `XA COMMIT` or `XA ROLLBACK` statement.
- If your MySQL installation contains a large amount of data that might take a long time to convert after an in-place upgrade, you might find it useful to create a “dummy” database instance for assessing what conversions might be needed and the work involved to perform them. Make a copy of your MySQL instance that contains a full copy of the `mysql` database, plus all other databases without data. Run your upgrade procedure on this dummy instance to see what actions might be needed so that you can better evaluate the work involved when performing actual data conversion on your original database instance.
- Rebuilding and reinstalling MySQL language interfaces is recommended whenever you install or upgrade to a new release of MySQL. This applies to MySQL interfaces such as PHP `mysql` extensions, the Perl `DBD::mysql` module, and the Python `MySQLdb` module.

Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation

Before starting the MySQL 8.0 server to upgrade the data directory, it is very important to ensure the upgrade readiness of your installation by using your MySQL 5.7 server to perform several preliminary checks. The upgrade process may fail otherwise.

To perform the preliminary checks:

1. There must be no tables that use obsolete data types, `InnoDB` tables that use nonnative partitioning, or triggers that have a missing or empty definer or an invalid creation context. (This context is indicated by the `character_set_client`, `collation_connection`, `Database Collation` attributes displayed by `SHOW TRIGGERS` or the `INFORMATION_SCHEMA TRIGGERS` table.) To identify tables and triggers that fail these requirements, execute this command:

   ```
   mysqlcheck -u root -p --all-databases --check-upgrade
   ```

   If `mysqlcheck` reports any errors, correct the issues.

2. There must be no partitioned tables that use a storage engine that does not have native partitioning support. To identify such tables, execute this query:

   ```
   SELECT TABLE_SCHEMA, TABLE_NAME
   FROM INFORMATION_SCHEMA.TABLES
   WHERE ENGINE NOT IN ('innodb', 'ndbcluster')
   AND CREATE_OPTIONS LIKE '%partitioned%';
   ```

   Any table reported by the query must be altered to use `InnoDB` or be made nonpartitioned. To change a table storage engine to `InnoDB`, execute this statement:

   ```
   ALTER TABLE table_name ENGINE = INNODB;
   ```

   Note

   For information about converting `MyISAM` tables to `InnoDB`, see Section 15.8.1.4, “Converting Tables from MyISAM to InnoDB”.

   To make a partitioned table nonpartitioned, execute this statement:

   ```
   ALTER TABLE table_name REMOVE PARTITIONING;
   ```

3. There must be no tables in the MySQL 5.7 `mysql` system database that have the same name as a table used by the MySQL 8.0 data dictionary. To identify tables with those names, execute this query:

```
SELECT TABLE_SCHEMA, TABLE_NAME
FROM INFORMATION_SCHEMA.TABLES
WHERE LOWER(TABLE_SCHEMA) = 'mysql'
and LOWER(TABLE_NAME) IN
(
'catalogs',
'character_sets',
'collations',
'column_type_elements',
'columns',
'events',
'foreign_key_column_usage',
'foreign_keys',
'index_column_usage',
'index_partitions',
'index_stats',
'indexes',
'parameter_type_elements',
'parameters',
'routines',
'schemata',
'st_spatial_reference_systems',
'table_partition_values',
'table_partitions',
'table_stats',
'tables',
'tablespace_files',
'tablespaces',
'triggers',
'version',
'view_routine_usage',
'view_table_usage'
);
```

Note

As of MySQL 8.0.2, replace `'version'` in the previous query with `'dd_properties'`.

Any tables reported by the query must be renamed (use `RENAME TABLE`). This may also entail changes to applications that use the affected tables.

4. There must be no tables that have foreign key constraint names longer than 64 characters. To identify tables with too-long constraint names, execute this query:

```
SELECT CONSTRAINT_SCHEMA, TABLE_NAME, CONSTRAINT_NAME
FROM INFORMATION_SCHEMA.REFERENTIAL_CONSTRAINTS
WHERE LENGTH(CONSTRAINT_NAME) > 64;
```

Any tables reported by the query must be altered to have constraint names no longer than 64 characters (use `ALTER TABLE`).

5. Your MySQL 5.7 installation must not use features that are not supported by MySQL 8.0. Any changes here are necessarily installation specific, but the following examples illustrate the kind of things to look for:

   - Tables that use a storage engine not supported in MySQL 8.0 must be altered to use a supported engine.

     Example: MySQL 8.0 does not yet support MySQL Cluster, so `NDB` tables must be altered to use a different storage engine.

   - Some server startup options and system variables have been removed in MySQL 8.0. See Features Removed in MySQL 8.0, and Section 1.5, “Server and Status Variables and Options

Added, Deprecated, or Removed in MySQL 8.0". If you use any of these items, an upgrade requires configuration changes.

Example: Because the data dictionary provides information about database objects, the server no longer checks directory names in the data directory to find databases. Consequently, the `--ignore-db-dir` option is extraneous and has been removed. To handle this, remove any instances of `--ignore-db-dir` from your startup configuration. In addition, remove or move the named data directory subdirectories before upgrading to MySQL 8.0. (Alternatively, let the 8.0 server add those directories to the data dictionary as databases, then remove each of those databases using `DROP DATABASE`.)

Performing an In-Place Upgrade

This section describes how to perform an in-place upgrade. Before proceeding, review Before You Begin, and ensure that your installation is upgrade-ready by following the procedure given in Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation.

Note

If you upgrade an installation originally produced by installing multiple RPM packages, upgrade all the packages, not just some. For example, if you previously installed the server and client RPMs, do not upgrade just the server RPM.

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, `mysqld_safe` is not installed. In such cases, use systemd for server startup and shutdown instead of the methods used in the following instructions. See Section 2.5.9, "Managing MySQL Server with systemd".

To perform an in-place upgrade:

1. Review the changes described in Section 2.10.1.1, "Changes Affecting Upgrades to MySQL 8.0" for steps to be performed before upgrading.
2. With your MySQL 5.7 server, if there are encrypted `InnoDB` tablespaces, rotate the keyring master key by executing this statement:

   ```
   ALTER INSTANCE ROTATE INNODB MASTER KEY;
   ```

3. Configure your MySQL 5.7 server to perform a slow shutdown by executing this statement:

   ```
   SET GLOBAL innodb_fast_shutdown = 0;
   ```

   With a slow shutdown, `InnoDB` performs a full purge and change buffer merge before shutting down, which ensures that the undo logs are empty and data files are fully prepared in case of file format differences between releases.

4. Shut down the old MySQL server. For example:

   ```
   mysqladmin -u root -p shutdown
   ```

5. Upgrade the MySQL binary installation or packages. If upgrading a binary installation, unpack the new MySQL binary distribution package. See Obtain and Unpack the Distribution. For package-based installations, replace the old packages with the new ones.

   Note

   For supported Linux distributions, the preferred method for replacing the MySQL packages is to use the MySQL software repositories; see

Section 2.10.1.2, “Upgrading MySQL with the MySQL Yum Repository”, Section 2.10.1.3, “Upgrading MySQL with the MySQL APT Repository”, or Upgrading MySQL with the MySQL SLES Repository for instructions.

6. Start the MySQL 8.0 server, using the existing data directory. For example:

```
mysqld_safe --user=mysql --datadir=/path/to/existing-datadir
```

If there are encrypted `InnoDB` tablespaces, use the `--early-plugin-load` option to load the keyring plugin.

The server automatically detects whether data dictionary tables are present. If not, the server creates them in the data directory, populates them with metadata, and then proceeds with its normal startup sequence. During this process, the server upgrades metadata for all database objects, including databases, tablespaces, system and user tables, views, and stored programs (stored procedures and functions, triggers, Event Scheduler events). The server also removes files that previously were used for metadata storage. For example, after upgrading, you will notice that your tables no longer have `.frm` files.

If this step succeeds, the server performs a cleanup:

- In the data directory, the server creates a directory named `backup_metadata_57` and moves into it files named `db.opt` and files with a suffix of `.frm`, `.par`, `.TRG`, `.TRN`, or `.isl`. (These are files previously used for metadata storage.)

  Files in the `backup_metadata_57` directory retain the original file system hierarchy. For example, if `t1.frm` was located in the `my_schema1` directory under the data directory, the server moves it to the `backup_metadata_57/my_schema1` directory.

- In the `mysql` database, the server renames the `event` and `proc` tables to `event_backup_57` and `proc_backup_57`.

If this step fails, the server reverts all changes to the data directory. In this case, you should remove all redo log files, start your MySQL 5.7 server on the same data directory, and fix the cause of any errors. Then perform another slow shutdown of the 5.7 server and start the MySQL 8.0 server to try again.

7. After the MySQL 8.0 server starts successfully, execute `mysql_upgrade`:

```
mysql_upgrade -u root -p
```

`mysql_upgrade` examines all tables in all databases for incompatibilities with the current version of MySQL. It makes any remaining changes required in the `mysql` system database between MySQL 5.7 and MySQL 8.0, so that you can take advantage of new privileges or capabilities. `mysql_upgrade` also brings the Performance Schema, `INFORMATION_SCHEMA`, and `sys` schema objects up to date for MySQL 8.0.

Note

`mysql_upgrade` should not be used when the server is running with `--gtid-mode=ON`. See GTID mode and mysql_upgrade for more information.

`mysql_upgrade` does not upgrade the contents of the help tables. For upgrade instructions, see Section 5.1.10, “Server-Side Help”.

8. Shut down and restart the MySQL server to ensure that any changes made to the system tables take effect. For example:

```
mysqladmin -u root -p shutdown
```

```
mysqld_safe --user=mysql --datadir=/path/to/existing-datadir
```

The first time you started the MySQL 8.0 server (in an earlier step), you may have noticed that it wrote several messages to its error log regarding nonupgraded tables. If `mysql_upgrade` has been run successfully, there should be no such messages the second time you start the server.

Performing a Logical Upgrade

This section describes how to perform a logical upgrade. Before proceeding, review Before You Begin, and ensure that your installation is upgrade-ready by following the procedure given in Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, `mysqld_safe` is not installed. In such cases, use systemd for server startup and shutdown instead of the methods used in the following instructions. See Section 2.5.9, “Managing MySQL Server with systemd”.

To perform a logical upgrade:

1. Review the changes described in Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0” for steps to be performed before upgrading.

2. With your MySQL 5.7 server, use `mysqldump` to export your data:

   ```
   mysqldump -u root -p
     --add-drop-table --routines --events --all-databases
     --ignore-table=mysql.innodb_table_stats
     --ignore-table=mysql.innodb_index_stats
     --force > data-for-upgrade.sql
   ```

   The `--routines` and `--events` options must be given explicitly to include stored routines and events as object definitions in the dump file. Otherwise, they will appear only as rows in the `mysql.proc` and `mysql.event` tables, which will not be used in MySQL 8.0 when the dump file is reloaded (definitions for the corresponding objects are expected to be stored in data dictionary tables instead).

   You can omit the `--ignore-table` options; the dump file will still reload later if you include the `--force` option on the load command. However, in that case, be prepared to see several warnings when those two tables load, which you should ignore.

3. Shut down the old MySQL server. For example:

   ```
   mysqladmin -u root -p shutdown
   ```

4. Install MySQL 8.0. For installation instructions, see Chapter 2, *Installing and Upgrading MySQL*.

5. Initialize a new data directory, as described at Section 2.9.1, “Initializing the Data Directory”. For example:

   ```
   mysqld --initialize --datadir=/path/to/8.0-datadir
   ```

   Copy the temporary `'root'@'localhost'` password displayed to your screen or written to your error log for later use.

6. Start the MySQL 8.0 server, using the new data directory. For example:

```
mysqld_safe --user=mysql --datadir=/path/to/8.0-datadir
```

If there are encrypted `InnoDB` tablespaces, use the `--early-plugin-load` option to load the keyring plugin.

7. Reset the temporary `root` password:

```
shell> mysql -u root -p
Enter password: ****  <- enter temporary root password
mysql> ALTER USER USER() IDENTIFIED BY 'your new password';
```

8. Load the previously created dump file into the new MySQL server. For example:

```
mysql -u root -p --force < data-for-upgrade.sql
```

If you did not specify `--ignore-table` options to suppress dumping the `mysql.event` and `mysql.proc` tables for the earlier `mysqldump` command, they will reappear when you reload the dump file.

9. After loading the dump file, execute `mysql_upgrade`:

```
mysql_upgrade -u root -p
```

`mysql_upgrade` examines all tables in all databases for incompatibilities with the current version of MySQL. It makes any remaining changes required in the `mysql` system database between MySQL 5.7 and MySQL 8.0, so that you can take advantage of new privileges or capabilities. `mysql_upgrade` also brings the Performance Schema, `INFORMATION_SCHEMA`, and `sys` schema objects up to date for MySQL 8.0.

Note

`mysql_upgrade` should not be used when the server is running with `--gtid-mode=ON`. See GTID mode and mysql_upgrade for more information.

`mysql_upgrade` does not upgrade the contents of the help tables. For upgrade instructions, see Section 5.1.10, “Server-Side Help”.

10. Shut down and restart the MySQL server to ensure that any changes made to the system tables take effect. For example:

```
mysqladmin -u root -p shutdown
mysqld_safe --user=mysql --datadir=/path/to/8.0-datadir
```

Upgrade Troubleshooting

- If problems occur, such as that the new `mysqld` server does not start, verify that you do not have an old `my.cnf` file from your previous installation. You can check this with the `--print-defaults` option (for example, `mysqld --print-defaults`). If this command displays anything other than the program name, you have an active `my.cnf` file that affects server or client operation.
- If, after an upgrade, you experience problems with compiled client programs, such as `Commands out of sync` or unexpected core dumps, you probably have used old header or library files when compiling your programs. In this case, check the date for your `mysql.h` file and `libmysqlclient.a` library to verify that they are from the new MySQL distribution. If not, recompile your programs with the new headers and libraries. Recompilation might also be necessary for programs compiled against the shared client library if the library major version number has changed (for example, from `libmysqlclient.so.15` to `libmysqlclient.so.16`).
- If you have created a user-defined function (UDF) with a given name and upgrade MySQL to a version that implements a new built-in function with the same name, the UDF becomes inaccessible.

To correct this, use `DROP FUNCTION` to drop the UDF, and then use `CREATE FUNCTION` to re-create the UDF with a different nonconflicting name. The same is true if the new version of MySQL implements a built-in function with the same name as an existing stored function. See Section 9.2.4, “Function Name Parsing and Resolution”, for the rules describing how the server interprets references to different kinds of functions.

2.10.1.1 Changes Affecting Upgrades to MySQL 8.0

Before upgrading to MySQL 8.0, review the changes described in this section to identify upgrade issues that apply to your current MySQL installation and applications.

Note

In addition to the changes outlined in this section, review the Release Notes and other important information outlined in Before You Begin.

Changes marked as either **Known issue** or **Incompatible change** are incompatibilities with earlier versions of MySQL, and may require your attention *before you upgrade*. Our aim is to avoid these changes, but occasionally they are necessary to correct problems that would be worse than an incompatibility between releases. If any upgrade issue applicable to your installation involves an incompatibility that requires special handling, follow the instructions given in the incompatibility description. Sometimes this involves dumping and reloading tables, or use of a statement such as `CHECK TABLE` or `REPAIR TABLE`.

For dump and reload instructions, see Section 2.10.3, “Rebuilding or Repairing Tables or Indexes”. Any procedure that involves `REPAIR TABLE` with the `USE_FRM` option *must* be done before upgrading. Use of this statement with a version of MySQL different from the one used to create the table (that is, using it after upgrading) may damage the table. See Section 13.7.2.5, “REPAIR TABLE Syntax”.

- Data Dictionary Changes
- Configuration Changes
- InnoDB Changes

Data Dictionary Changes

Incompatible change: MySQL Server 8.0 incorporates a global data dictionary containing information about database objects in transactional tables. In previous MySQL series, dictionary data was stored in metadata files and nontransactional system tables. As a result, the upgrade procedure is somewhat different from previous MySQL releases and requires that you verify the upgrade readiness of your installation by checking specific prerequisites. For more information, see Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation. A data dictionary-enabled server entails some general operational differences; see Section 14.6, “Data Dictionary Usage Differences”.

Configuration Changes

- **Incompatible change**: A MySQL storage engine is now responsible for providing its own partitioning handler, and the MySQL server no longer provides generic partitioning support. `InnoDB` is the only storage engine providing a native partitioning handler that is supported in MySQL 8.0. (The `NDB` storage engine also provides native partitioning support, but it is not yet supported in MySQL 8.0.) A partitioned table using any other storage engine must be altered—either to convert it to `InnoDB`, or to remove its partitioning—*before* upgrading the server, else it cannot be used afterwards.

  For information about converting `MyISAM` tables to `InnoDB`, see Section 15.8.1.4, “Converting Tables from MyISAM to InnoDB”.

  A table creation statement that would result in a partitioned table using a storage engine without such support fails with an error (`ER_CHECK_NOT_IMPLEMENTED`) in MySQL 8.0. If you import databases from a dump file created in MySQL 5.7 (or earlier) using `mysqldump` into a MySQL 8.0 server, you

must make sure that any statements creating partitioned tables do not also specify an unsupported storage engine, either by removing any references to partitioning, or by specifying the storage engine as `InnoDB` or allowing it to be set as `InnoDB` by default.

Note

The procedure given at Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation, describes how to identify partitioned tables that must be altered before upgrading to MySQL 8.0.

See Section 22.6.2, “Partitioning Limitations Relating to Storage Engines”, for further information.

- **Incompatible change**: Several server error codes are not used and have been removed (for a list, see Features Removed in MySQL 8.0). Applications that test specifically for any of them should be updated.

InnoDB Changes

- **Incompatible change**: Due to redo log format differences, performing an in-place upgrade to MySQL 8.0 from MySQL 5.7 requires that MySQL 5.7 be shut down using an `innodb_fast_shutdown` setting of `0` before upgrading. (It is required that you upgrade from a GA version of MySQL 5.7; that is, 5.7.9 or higher. Upgrades from non-GA versions of MySQL 5.7 are not supported.) Performing a shutdown using `innodb_fast_shutdown=0` is a recommended step in Performing an In-Place Upgrade.

SQL Changes

- Some keywords may be reserved in MySQL 8.0 that were not reserved in MySQL 5.7. See Section 9.3, “Keywords and Reserved Words”. This can cause words previously used as identifiers to become illegal. To fix affected statements, use identifier quoting. See Section 9.2, “Schema Object Names”.
- After upgrading, it is recommended that you test optimizer hints specified in application code to ensure that the hints are still required to achieve the desired optimization strategy. Optimizer enhancements can sometimes render certain optimizer hints unnecessary. In some cases, an unnecessary optimizer hint may even be counterproductive.

2.10.1.2 Upgrading MySQL with the MySQL Yum Repository

For supported Yum-based platforms (see Section 2.5.1, “Installing MySQL on Linux Using the MySQL Yum Repository”, for a list), you can perform an in-place upgrade for MySQL (that is, replacing the old version and then running the new version off the old data files) with the MySQL Yum repository.

Notes

- Before performing any update to MySQL, follow carefully the instructions in Section 2.10.1, “Upgrading MySQL”. Among other instructions discussed there, it is especially important to back up your database before the update.
- The following instructions assume you have installed MySQL with the MySQL Yum repository or with an RPM package directly downloaded from MySQL Developer Zone's MySQL Download page; if that is not the case, following the instructions in Replacing a Third-Party Distribution of MySQL Using the MySQL Yum Repository.

Selecting a Target Series

By default, the MySQL Yum repository updates MySQL to the latest version in the release series you have chosen during installation (see Selecting a Release Series for details), which means, for example, a 5.7.x installation will NOT be updated to a 8.0.x release automatically. To update

to another release series, you need to first disable the subrepository for the series that has been selected (by default, or by yourself) and enable the subrepository for your target series. To do that, see the general instructions given in Selecting a Release Series. For upgrading from MySQL 5.7 to 8.0, perform the *reverse* of the steps illustrated in Selecting a Release Series, disabling the subrepository for the MySQL 5.7 series and enabling that for the MySQL 8.0 series.

As a general rule, to upgrade from one release series to another, go to the next series rather than skipping a series. For example, if you are currently running MySQL 5.6 and wish to upgrade to 8.0, upgrade to MySQL 5.7 first before upgrading to 8.0.

Important

For important information about upgrading from MySQL 5.7 to 8.0, see Upgrading from MySQL 5.7 to 8.0.

2. **Upgrading MySQL**

Upgrade MySQL and its components by the following command, for platforms that are not dnf-enabled:

```
sudo yum update mysql-server
```

For platforms that are dnf-enabled:

```
sudo dnf upgrade mysql-server
```

Alternatively, you can update MySQL by telling Yum to update everything on your system, which might take considerably more time; for platforms that are not dnf-enabled:

```
sudo yum update
```

For platforms that are dnf-enabled:

```
sudo dnf upgrade
```

3. **Restarting MySQL**

The MySQL server always restarts after an update by Yum. Once the server restarts, run `mysql_upgrade` to check and possibly resolve any incompatibilities between the old data and the upgraded software. `mysql_upgrade` also performs other functions; see Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables” for details.

You can also update only a specific component. Use the following command to list all the installed packages for the MySQL components (for dnf-enabled systems, replace `yum` in the command with `dnf`):

```
sudo yum list installed | grep "^mysql"
```

After identifying the package name of the component of your choice, for platforms that are not dnf-enabled, update the package with the following command, replacing *`package-name`* with the name of the package:

```
sudo yum update package-name
```

For dnf-enabled platforms:

```
sudo dnf upgrade package-name
```

Upgrading the Shared Client Libraries

After updating MySQL using the Yum repository, applications compiled with older versions of the shared client libraries should continue to work.

If you recompile applications and dynamically link them with the updated libraries: As typical with new versions of shared libraries where there are differences or additions in symbol versioning between the newer and older libraries (for example, between the newer, standard 8.0 shared client libraries and some older—prior or variant—versions of the shared libraries shipped natively by the Linux distributions' software repositories, or from some other sources), any applications compiled using the updated, newer shared libraries will require those updated libraries on systems where the applications are deployed. And, as expected, if those libraries are not in place, the applications requiring the shared libraries will fail. So, be sure to deploy the packages for the shared libraries from MySQL on those systems. You can do this by adding the MySQL Yum repository to the systems (see Adding the MySQL Yum Repository) and install the latest shared libraries using the instructions given in Installing Additional MySQL Products and Components with Yum.

2.10.1.3 Upgrading MySQL with the MySQL APT Repository

On Debian and Ubuntu platforms, you can perform an in-place upgrade of MySQL and its components with the MySQL APT repository. See Upgrading MySQL with the MySQL APT Repository in A Quick Guide to Using the MySQL APT Repository.

2.10.2 Downgrading MySQL

This section describes how to downgrade to an older MySQL version.

- Supported Downgrade Methods
- Supported Downgrade Paths
- Before You Begin
- Performing an In-Place Downgrade
- Performing a Logical Downgrade
- Downgrade Troubleshooting

Note

In the following discussion, MySQL commands that must be run using a MySQL account with administrative privileges include `-u root` on the command line to specify the MySQL `root` user. Commands that require a password for `root` also include a `-p` option. Because `-p` is followed by no option value, such commands prompt for the password. Type the password when prompted and press Enter.

SQL statements can be executed using the `mysql` command-line client (connect as `root` to ensure that you have the necessary privileges).

Supported Downgrade Methods

Supported downgrade methods include:

- *In-Place Downgrade*: Involves shutting down the new MySQL version, replacing the new MySQL binaries or packages with the old ones, and restarting the old MySQL version on the existing data directory. In-place downgrades are supported for downgrades between GA versions within the same release series.
- *Logical Downgrade*: Involves using `mysqldump` to dump all tables from the new MySQL version, and then loading the dump file into the old MySQL version. Logical downgrades are supported for

downgrades between GA versions within the same release series and for downgrades between release levels.

Note

In-place downgrades are not supported for MySQL APT, SLES, and Yum repository installations.

For procedures, see Performing an In-Place Downgrade, and Performing a Logical Downgrade.

Supported Downgrade Paths

Unless otherwise documented, the following downgrade paths are supported:

- Downgrading from a release series version to an older release series version is supported using all downgrade methods.
- Downgrading one release level is supported using the *logical downgrade* method. For example, downgrading from 8.0 to 5.7 is supported.
- Downgrading more than one release level is supported using the *logical downgrade* method, but only if you downgrade one release level at a time. For example, you can downgrade from 8.0 to 5.7, and then to 5.6.

The following conditions apply to all downgrade paths:

- Downgrades between General Availability (GA) status releases are supported.
- Downgrades between milestone releases (or from a GA release to a milestone release) are not supported. For example, downgrading from MySQL 8.0.1 to MySQL 8.0.0 is not supported, as neither are GA status releases.

Before You Begin

Before downgrading, the following steps are recommended:

- Review the Release Notes for the MySQL version you are downgrading from to ensure that there are no features or fixes that you really need.
- Review Section 2.10.2.1, “Changes Affecting Downgrades from MySQL 8.0”. This section describes changes that may require action before or after downgrading.

Note

The downgrade procedures described in the following sections assume you are downgrading with data files created or modified by the newer MySQL version. However, if you did not modify your data after upgrading, downgrading using backups taken *before* upgrading to the new MySQL version is recommended. Many of the changes described in Section 2.10.2.1, “Changes Affecting Downgrades from MySQL 8.0” that require action before or after downgrading are not applicable when downgrading using backups taken *before* upgrading to the new MySQL version.

- Always back up your current databases and log files before downgrading. The backup should include the `mysql` database, which contains the MySQL system tables. See Section 7.2, “Database Backup Methods”.
- Use of new features, new configuration options, or new configuration option values that are not supported by a previous release may cause downgrade errors or failures. Before downgrading, it is recommended that you reverse changes resulting from the use of new features and remove configuration settings that are not supported by the release you are downgrading to.

- If you use XA transactions with `InnoDB`, run `XA RECOVER` before downgrading to check for uncommitted XA transactions. If results are returned, either commit or rollback the XA transactions by issuing an `XA COMMIT` or `XA ROLLBACK` statement.

Performing an In-Place Downgrade

In-place downgrades are supported for downgrades between GA status releases within the same release series. Before proceeding, review Before You Begin.

Warning

For a supported downgrade path within the MySQL 8.0 release series, there must be at least two MySQL 8.0 GA status versions available.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, `mysqld_safe` is not installed. In such cases, use systemd for server startup and shutdown instead of the methods used in the following instructions. See Section 2.5.9, “Managing MySQL Server with systemd”.

To perform an in-place downgrade:

1. Review the changes described in Section 2.10.2.1, “Changes Affecting Downgrades from MySQL 8.0” for steps to be performed before downgrading.

2. Configure MySQL to perform a slow shutdown by setting `innodb_fast_shutdown` to `0`. For example:

   ```
   mysql -u root -p --execute="SET GLOBAL innodb_fast_shutdown=0"
   ```

   With a slow shutdown, `InnoDB` performs a full purge and change buffer merge before shutting down, which ensures that data files are fully prepared in case of file format differences between releases.

3. Shut down the newer MySQL server. For example:

   ```
   mysqladmin -u root -p shutdown
   ```

4. After the slow shutdown, remove the `InnoDB` redo log files (the `ib_logfile*` files) from the `data` directory to avoid downgrade issues related to redo log file format changes that may have occurred between releases.

   ```
   rm ib_logfile*
   ```

5. Downgrade the MySQL binaries or packages in-place by replacing the newer binaries or packages with the older ones.

6. Start the older (downgraded) MySQL server, using the existing data directory. For example:

   ```
   mysqld_safe --user=mysql --datadir=/path/to/existing-datadir
   ```

7. Run `mysql_upgrade`. For example:

   ```
   mysql_upgrade -u root -p
   ```

   `mysql_upgrade` examines all tables in all databases for incompatibilities with the current version of MySQL, and attempts to repair the tables if problems are found.

8. Shut down and restart the MySQL server to ensure that any changes made to the system tables take effect. For example:

```
mysqladmin -u root -p shutdown
mysqld_safe --user=mysql --datadir=/path/to/existing-datadir
```

Performing a Logical Downgrade

Logical downgrades are supported for downgrades between releases within the same release series and for downgrades to the previous release level. Only downgrades between General Availability (GA) status releases are supported. Before proceeding, review Before You Begin.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, `mysqld_safe` is not installed. In such cases, use systemd for server startup and shutdown instead of the methods used in the following instructions. See Section 2.5.9, “Managing MySQL Server with systemd”.

For MySQL APT, SLES, and Yum repository installations, only downgrades to the previous release level are supported. Where the instructions call for initializing an older instance, use the package management utility to remove MySQL 8.0 packages and install MySQL 5.7 packages.

To perform a logical downgrade:

1. Review the changes described in Section 2.10.2.1, “Changes Affecting Downgrades from MySQL 8.0” for steps to be performed before downgrading.

2. Dump all databases. For example:

```
mysqldump -u root -p
  --add-drop-table --routines --events
  --all-databases --force > data-for-downgrade.sql
```

3. Shut down the newer MySQL server. For example:

```
mysqladmin -u root -p shutdown
```

4. Initialize an older MySQL instance, with a new data directory. For example:

```
mysqld --initialize --user=mysql
```

5. Start the older MySQL server, using the new data directory. For example:

```
mysqld_safe --user=mysql --datadir=/path/to/new-datadir
```

6. Load the dump file into the older MySQL server. For example:

```
mysql -u root -p --force < data-for-upgrade.sql
```

7. Run `mysql_upgrade`. For example:

```
mysql_upgrade -u root -p
```

`mysql_upgrade` examines all tables in all databases for incompatibilities with the current version of MySQL, and attempts to repair the tables if problems are found.

8. Shut down and restart the MySQL server to ensure that any changes made to the system tables take effect. For example:

```
mysqladmin -u root -p shutdown
mysqld_safe --user=mysql --datadir=/path/to/new-datadir
```

Downgrade Troubleshooting

If you downgrade from one release series to another, there may be incompatibilities in table storage formats. In this case, use `mysqldump` to dump your tables before downgrading. After downgrading, reload the dump file using `mysql` or `mysqlimport` to re-create your tables. For examples, see Section 2.10.4, “Copying MySQL Databases to Another Machine”.

A typical symptom of a downward-incompatible table format change when you downgrade is that you cannot open tables. In that case, use the following procedure:

1. Stop the older MySQL server that you are downgrading to.
2. Restart the newer MySQL server you are downgrading from.
3. Dump any tables that were inaccessible to the older server by using `mysqldump` to create a dump file.
4. Stop the newer MySQL server and restart the older one.
5. Reload the dump file into the older server. Your tables should be accessible.

2.10.2.1 Changes Affecting Downgrades from MySQL 8.0

Before downgrading from MySQL 8.0, review the changes described in this section. Some changes may require action before or after downgrading.

- System Table Changes
- InnoDB Changes

System Table Changes

- In MySQL 8.0.2, the system tables in the `mysql` system database were migrated from individual tablespace files to a single tablespace file (`mysql.ibd`). Before downgrading to MySQL 5.7, move the `mysql` system tables back to individual tablespace files by running the following statements:

```
ALTER TABLE mysql.columns_priv TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.component TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.db TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.default_roles TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.engine_cost TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.func TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.general_log TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.global_grants TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.gtid_executed TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.help_category TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.help_keyword TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.help_relation TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.help_topic TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.innodb_index_stats TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.innodb_table_stats TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.plugin TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.procs_priv TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.proxies_priv TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.role_edges TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.server_cost TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.servers TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.slave_master_info TABLESPACE=innodb_file_per_table;
```

```
ALTER TABLE mysql.slave_relay_log_info TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.slave_worker_info TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.slow_log TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.tables_priv TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.time_zone TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.time_zone_leap_second TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.time_zone_name TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.time_zone_transition TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.time_zone_transition_type TABLESPACE=innodb_file_per_table;
ALTER TABLE mysql.user TABLESPACE=innodb_file_per_table;
```

- In MySQL 8.0.2, six system tables changed from `MyISAM` to `InnoDB`. Before downgrading to MySQL 5.7, change each table back to `MyISAM` by running the following statements:

```
ALTER TABLE mysql.columns_priv ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT
ALTER TABLE mysql.db ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT;
ALTER TABLE mysql.procs_priv ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT;
ALTER TABLE mysql.proxies_priv ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT;
ALTER TABLE mysql.tables_priv ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT;
ALTER TABLE mysql.user ENGINE='MyISAM' STATS_PERSISTENT=DEFAULT;
```

- In MySQL 8.0.2, two columns were added to the `mysql.user` table. Before downgrading to MySQL 5.7, remove these columns by running the following statements:

```
ALTER TABLE mysql.user drop Create_role_priv;
ALTER TABLE mysql.user drop Drop_role_priv;
```

InnoDB Changes

- **Incompatible change**: Before performing an in-place downgrade from MySQL 8.0 to 5.7, you must shutdown MySQL using an `innodb_fast_shutdown` setting of `0` or `1` followed by removal of the redo log files (the `ib_logfile*` files). Performing a shutdown using `innodb_fast_shutdown=0` and removing the redo logs are recommended steps in Performing an In-Place Downgrade.

2.10.3 Rebuilding or Repairing Tables or Indexes

This section describes how to rebuild or repair tables or indexes, which may be necessitated by:

- Changes to how MySQL handles data types or character sets. For example, an error in a collation might have been corrected, necessitating a table rebuild to update the indexes for character columns that use the collation.
- Required table repairs or upgrades reported by `CHECK TABLE`, `mysqlcheck`, or `mysql_upgrade`.

Methods for rebuilding a table include:

- Dump and Reload Method
- ALTER TABLE Method
- REPAIR TABLE Method

Dump and Reload Method

If you are rebuilding tables because a different version of MySQL will not handle them after a binary (in-place) upgrade or downgrade, you must use the dump-and-reload method. Dump the tables *before* upgrading or downgrading using your original version of MySQL. Then reload the tables *after* upgrading or downgrading.

If you use the dump-and-reload method of rebuilding tables only for the purpose of rebuilding indexes, you can perform the dump either before or after upgrading or downgrading. Reloading still must be done afterward.

If you need to rebuild an `InnoDB` table because a `CHECK TABLE` operation indicates that a table upgrade is required, use `mysqldump` to create a dump file and `mysql` to reload the file. If the `CHECK TABLE` operation indicates that there is a corruption or causes `InnoDB` to fail, refer to Section 15.20.2, "Forcing InnoDB Recovery" for information about using the `innodb_force_recovery` option to restart `InnoDB`. To understand the type of problem that `CHECK TABLE` may be encountering, refer to the `InnoDB` notes in Section 13.7.2.2, "CHECK TABLE Syntax".

To rebuild a table by dumping and reloading it, use `mysqldump` to create a dump file and `mysql` to reload the file:

```
mysqldump db_name t1 > dump.sql
mysql db_name < dump.sql
```

To rebuild all the tables in a single database, specify the database name without any following table name:

```
mysqldump db_name > dump.sql
mysql db_name < dump.sql
```

To rebuild all tables in all databases, use the `--all-databases` option:

```
mysqldump --all-databases > dump.sql
mysql < dump.sql
```

ALTER TABLE Method

To rebuild a table with `ALTER TABLE`, use a "null" alteration; that is, an `ALTER TABLE` statement that "changes" the table to use the storage engine that it already has. For example, if `t1` is an `InnoDB` table, use this statement:

```
ALTER TABLE t1 ENGINE = InnoDB;
```

If you are not sure which storage engine to specify in the `ALTER TABLE` statement, use `SHOW CREATE TABLE` to display the table definition.

REPAIR TABLE Method

The `REPAIR TABLE` method is only applicable to `MyISAM`, `ARCHIVE`, and `CSV` tables.

You can use `REPAIR TABLE` if the table checking operation indicates that there is a corruption or that an upgrade is required. For example, to repair a `MyISAM` table, use this statement:

```
REPAIR TABLE t1;
```

`mysqlcheck --repair` provides command-line access to the `REPAIR TABLE` statement. This can be a more convenient means of repairing tables because you can use the `--databases` or `--all-databases` option to repair all tables in specific databases or all databases, respectively:

```
mysqlcheck --repair --databases db_name ...
mysqlcheck --repair --all-databases
```

2.10.4 Copying MySQL Databases to Another Machine

In cases where you need to transfer databases between different architectures, you can use `mysqldump` to create a file containing SQL statements. You can then transfer the file to the other machine and feed it as input to the `mysql` client.

Use `mysqldump --help` to see what options are available.

The easiest (although not the fastest) way to move a database between two machines is to run the following commands on the machine on which the database is located:

```
mysqladmin -h 'other_hostname' create db_name
mysqldump db_name | mysql -h 'other_hostname' db_name
```

If you want to copy a database from a remote machine over a slow network, you can use these commands:

```
mysqladmin create db_name
mysqldump -h 'other_hostname' --compress db_name | mysql db_name
```

You can also store the dump in a file, transfer the file to the target machine, and then load the file into the database there. For example, you can dump a database to a compressed file on the source machine like this:

```
mysqldump --quick db_name | gzip > db_name.gz
```

Transfer the file containing the database contents to the target machine and run these commands there:

```
mysqladmin create db_name
gunzip < db_name.gz | mysql db_name
```

You can also use `mysqldump` and `mysqlimport` to transfer the database. For large tables, this is much faster than simply using `mysqldump`. In the following commands, *`DUMPDIR`* represents the full path name of the directory you use to store the output from `mysqldump`.

First, create the directory for the output files and dump the database:

```
mkdir DUMPDIR
mysqldump --tab=DUMPDIR db_name
```

Then transfer the files in the *`DUMPDIR`* directory to some corresponding directory on the target machine and load the files into MySQL there:

```
mysqladmin create db_name           # create database
cat DUMPDIR/*.sql | mysql db_name   # create tables in database
mysqlimport db_name DUMPDIR/*.txt   # load data into tables
```

Do not forget to copy the `mysql` database because that is where the grant tables are stored. You might have to run commands as the MySQL `root` user on the new machine until you have the `mysql` database in place.

After you import the `mysql` database on the new machine, execute `mysqladmin flush-privileges` so that the server reloads the grant table information.

2.11 Perl Installation Notes

The Perl `DBI` module provides a generic interface for database access. You can write a `DBI` script that works with many different database engines without change. To use `DBI`, you must install the `DBI` module, as well as a DataBase Driver (DBD) module for each type of database server you want to access. For MySQL, this driver is the `DBD::mysql` module.

Note

Perl support is not included with MySQL distributions. You can obtain the necessary modules from http://search.cpan.org for Unix, or by using the

ActiveState `ppm` program on Windows. The following sections describe how to do this.

The `DBI`/`DBD` interface requires Perl 5.6.0, and 5.6.1 or later is preferred. DBI *does not work* if you have an older version of Perl. You should use `DBD::mysql` 4.009 or higher. Although earlier versions are available, they do not support the full functionality of MySQL 8.0.

2.11.1 Installing Perl on Unix

MySQL Perl support requires that you have installed MySQL client programming support (libraries and header files). Most installation methods install the necessary files. If you install MySQL from RPM files on Linux, be sure to install the developer RPM as well. The client programs are in the client RPM, but client programming support is in the developer RPM.

The files you need for Perl support can be obtained from the CPAN (Comprehensive Perl Archive Network) at http://search.cpan.org.

The easiest way to install Perl modules on Unix is to use the `CPAN` module. For example:

```
shell> perl -MCPAN -e shell
cpan> install DBI
cpan> install DBD::mysql
```

The `DBD::mysql` installation runs a number of tests. These tests attempt to connect to the local MySQL server using the default user name and password. (The default user name is your login name on Unix, and `ODBC` on Windows. The default password is “no password.”) If you cannot connect to the server with those values (for example, if your account has a password), the tests fail. You can use `force install DBD::mysql` to ignore the failed tests.

`DBI` requires the `Data::Dumper` module. It may be installed; if not, you should install it before installing `DBI`.

It is also possible to download the module distributions in the form of compressed `tar` archives and build the modules manually. For example, to unpack and build a DBI distribution, use a procedure such as this:

1. Unpack the distribution into the current directory:

   ```
   shell> gunzip < DBI-VERSION.tar.gz | tar xvf -
   ```

   This command creates a directory named `DBI-VERSION`.

2. Change location into the top-level directory of the unpacked distribution:

   ```
   shell> cd DBI-VERSION
   ```

3. Build the distribution and compile everything:

   ```
   shell> perl Makefile.PL
   shell> make
   shell> make test
   shell> make install
   ```

The `make test` command is important because it verifies that the module is working. Note that when you run that command during the `DBD::mysql` installation to exercise the interface code, the MySQL server must be running or the test fails.

It is a good idea to rebuild and reinstall the `DBD::mysql` distribution whenever you install a new release of MySQL. This ensures that the latest versions of the MySQL client libraries are installed correctly.

If you do not have access rights to install Perl modules in the system directory or if you want to install local Perl modules, the following reference may be useful: http://learn.perl.org/faq/perlfaq8.html#How-do-I-keep-my-own-module-library-directory-

2.11.2 Installing ActiveState Perl on Windows

On Windows, you should do the following to install the MySQL `DBD` module with ActiveState Perl:

1. Get ActiveState Perl from http://www.activestate.com/Products/ActivePerl/ and install it.
2. Open a console window.
3. If necessary, set the `HTTP_proxy` variable. For example, you might try a setting like this:

```
C:\> set HTTP_proxy=my.proxy.com:3128
```

4. Start the PPM program:

```
C:\> C:\perl\bin\ppm.pl
```

5. If you have not previously done so, install `DBI`:

```
ppm> install DBI
```

6. If this succeeds, run the following command:

```
ppm> install DBD-mysql
```

This procedure should work with ActiveState Perl 5.6 or higher.

If you cannot get the procedure to work, you should install the ODBC driver instead and connect to the MySQL server through ODBC:

```
use DBI;
$dbh= DBI->connect("DBI:ODBC:$dsn",$user,$password) ||
  die "Got error $DBI::errstr when connecting to $dsn\n";
```

2.11.3 Problems Using the Perl DBI/DBD Interface

If Perl reports that it cannot find the `../mysql/mysql.so` module, the problem is probably that Perl cannot locate the `libmysqlclient.so` shared library. You should be able to fix this problem by one of the following methods:

- Copy `libmysqlclient.so` to the directory where your other shared libraries are located (probably `/usr/lib` or `/lib`).
- Modify the `-L` options used to compile `DBD::mysql` to reflect the actual location of `libmysqlclient.so`.
- On Linux, you can add the path name of the directory where `libmysqlclient.so` is located to the `/etc/ld.so.conf` file.
- Add the path name of the directory where `libmysqlclient.so` is located to the `LD_RUN_PATH` environment variable. Some systems use `LD_LIBRARY_PATH` instead.

Note that you may also need to modify the `-L` options if there are other libraries that the linker fails to find. For example, if the linker cannot find `libc` because it is in `/lib` and the link command specifies `-L/usr/lib`, change the `-L` option to `-L/lib` or add `-L/lib` to the existing link command.

If you get the following errors from `DBD::mysql`, you are probably using `gcc` (or using an old binary compiled with `gcc`):

```
/usr/bin/perl: can't resolve symbol '__moddi3'
/usr/bin/perl: can't resolve symbol '__divdi3'
```

Add `-L/usr/lib/gcc-lib/... -lgcc` to the link command when the `mysql.so` library gets built (check the output from `make` for `mysql.so` when you compile the Perl client). The `-L` option should specify the path name of the directory where `libgcc.a` is located on your system.

Another cause of this problem may be that Perl and MySQL are not both compiled with `gcc`. In this case, you can solve the mismatch by compiling both with `gcc`.

Chapter 3 Tutorial

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This chapter provides a tutorial introduction to MySQL by showing how to use the `mysql` client program to create and use a simple database. `mysql` (sometimes referred to as the “terminal monitor” or just “monitor”) is an interactive program that enables you to connect to a MySQL server, run queries, and view the results. `mysql` may also be used in batch mode: you place your queries in a file beforehand, then tell `mysql` to execute the contents of the file. Both ways of using `mysql` are covered here.

To see a list of options provided by `mysql`, invoke it with the `--help` option:

```
shell> mysql --help
```

This chapter assumes that `mysql` is installed on your machine and that a MySQL server is available to which you can connect. If this is not true, contact your MySQL administrator. (If *you* are the administrator, you need to consult the relevant portions of this manual, such as Chapter 5, *MySQL Server Administration*.)

This chapter describes the entire process of setting up and using a database. If you are interested only in accessing an existing database, you may want to skip over the sections that describe how to create the database and the tables it contains.

Because this chapter is tutorial in nature, many details are necessarily omitted. Consult the relevant sections of the manual for more information on the topics covered here.

3.1 Connecting to and Disconnecting from the Server

To connect to the server, you will usually need to provide a MySQL user name when you invoke `mysql` and, most likely, a password. If the server runs on a machine other than the one where you log in, you will also need to specify a host name. Contact your administrator to find out what connection parameters you should use to connect (that is, what host, user name, and password to use). Once you know the proper parameters, you should be able to connect like this:

```
shell> mysql -h host -u user -p
Enter password: ********
```

host and *user* represent the host name where your MySQL server is running and the user name of your MySQL account. Substitute appropriate values for your setup. The `********` represents your password; enter it when `mysql` displays the `Enter password:` prompt.

If that works, you should see some introductory information followed by a `mysql>` prompt:

```
shell> mysql -h host -u user -p
Enter password: ********
Welcome to the MySQL monitor.  Commands end with ; or \g.
Your MySQL connection id is 25338 to server version: 8.0.3-standard

Type 'help;' or '\h' for help. Type '\c' to clear the buffer.

mysql>
```

The `mysql>` prompt tells you that `mysql` is ready for you to enter SQL statements.

If you are logging in on the same machine that MySQL is running on, you can omit the host, and simply use the following:

```
shell> mysql -u user -p
```

If, when you attempt to log in, you get an error message such as `ERROR 2002 (HY000): Can't connect to local MySQL server through socket '/tmp/mysql.sock' (2)`, it means that the MySQL server daemon (Unix) or service (Windows) is not running. Consult the administrator or see the section of Chapter 2, *Installing and Upgrading MySQL* that is appropriate to your operating system.

For help with other problems often encountered when trying to log in, see Section B.5.2, “Common Errors When Using MySQL Programs”.

Some MySQL installations permit users to connect as the anonymous (unnamed) user to the server running on the local host. If this is the case on your machine, you should be able to connect to that server by invoking `mysql` without any options:

```
shell> mysql
```

After you have connected successfully, you can disconnect any time by typing `QUIT` (or `\q`) at the `mysql>` prompt:

```
mysql> QUIT
Bye
```

On Unix, you can also disconnect by pressing Control+D.

Most examples in the following sections assume that you are connected to the server. They indicate this by the `mysql>` prompt.

3.2 Entering Queries

Make sure that you are connected to the server, as discussed in the previous section. Doing so does not in itself select any database to work with, but that is okay. At this point, it is more important to find out a little about how to issue queries than to jump right in creating tables, loading data into them, and retrieving data from them. This section describes the basic principles of entering queries, using several queries you can try out to familiarize yourself with how `mysql` works.

Here is a simple query that asks the server to tell you its version number and the current date. Type it in as shown here following the `mysql>` prompt and press Enter:

```
mysql> SELECT VERSION(), CURRENT_DATE;
+-----------+--------------+
| VERSION() | CURRENT_DATE |
+-----------+--------------+
| 5.8.0-m17 | 2015-12-21   |
+-----------+--------------+
1 row in set (0.02 sec)
mysql>
```

This query illustrates several things about `mysql`:

- A query normally consists of an SQL statement followed by a semicolon. (There are some exceptions where a semicolon may be omitted. `QUIT`, mentioned earlier, is one of them. We'll get to others later.)
- When you issue a query, `mysql` sends it to the server for execution and displays the results, then prints another `mysql>` prompt to indicate that it is ready for another query.
- `mysql` displays query output in tabular form (rows and columns). The first row contains labels for the columns. The rows following are the query results. Normally, column labels are the names of the columns you fetch from database tables. If you're retrieving the value of an expression rather than a table column (as in the example just shown), `mysql` labels the column using the expression itself.
- `mysql` shows how many rows were returned and how long the query took to execute, which gives you a rough idea of server performance. These values are imprecise because they represent wall clock time (not CPU or machine time), and because they are affected by factors such as server load and network latency. (For brevity, the “rows in set” line is sometimes not shown in the remaining examples in this chapter.)

Keywords may be entered in any lettercase. The following queries are equivalent:

```
mysql> SELECT VERSION(), CURRENT_DATE;
mysql> select version(), current_date;
mysql> SeLeCt vErSiOn(), current_DATE;
```

Here is another query. It demonstrates that you can use `mysql` as a simple calculator:

```
mysql> SELECT SIN(PI()/4), (4+1)*5;
+------------------+---------+
| SIN(PI()/4)      | (4+1)*5 |
+------------------+---------+
| 0.70710678118655 |      25 |
+------------------+---------+
1 row in set (0.02 sec)
```

The queries shown thus far have been relatively short, single-line statements. You can even enter multiple statements on a single line. Just end each one with a semicolon:

```
mysql> SELECT VERSION(); SELECT NOW();
+-----------+
| VERSION() |
+-----------+
| 5.8.0-m17 |
+-----------+
1 row in set (0.00 sec)

+---------------------+
| NOW()               |
+---------------------+
| 2015-12-21 14:06:10 |
+---------------------+
```

```
1 row in set (0.00 sec)
```

A query need not be given all on a single line, so lengthy queries that require several lines are not a problem. `mysql` determines where your statement ends by looking for the terminating semicolon, not by looking for the end of the input line. (In other words, `mysql` accepts free-format input: it collects input lines but does not execute them until it sees the semicolon.)

Here is a simple multiple-line statement:

```
mysql> SELECT
    -> USER()
    -> ,
    -> CURRENT_DATE;
+---------------+--------------+
| USER()        | CURRENT_DATE |
+---------------+--------------+
| jon@localhost | 2015-12-21   |
+---------------+--------------+
```

In this example, notice how the prompt changes from `mysql>` to `->` after you enter the first line of a multiple-line query. This is how `mysql` indicates that it has not yet seen a complete statement and is waiting for the rest. The prompt is your friend, because it provides valuable feedback. If you use that feedback, you can always be aware of what `mysql` is waiting for.

If you decide you do not want to execute a query that you are in the process of entering, cancel it by typing `\c`:

```
mysql> SELECT
    -> USER()
    -> \c
mysql>
```

Here, too, notice the prompt. It switches back to `mysql>` after you type `\c`, providing feedback to indicate that `mysql` is ready for a new query.

The following table shows each of the prompts you may see and summarizes what they mean about the state that `mysql` is in.

Prompt	Meaning
`mysql>`	Ready for new query
`->`	Waiting for next line of multiple-line query
`'>`	Waiting for next line, waiting for completion of a string that began with a single quote (`'`)
`">`	Waiting for next line, waiting for completion of a string that began with a double quote (`"`)
`` `> ``	Waiting for next line, waiting for completion of an identifier that began with a backtick (`` ` ``)
`/*>`	Waiting for next line, waiting for completion of a comment that began with `/*`

Multiple-line statements commonly occur by accident when you intend to issue a query on a single line, but forget the terminating semicolon. In this case, `mysql` waits for more input:

```
mysql> SELECT USER()
    ->
```

If this happens to you (you think you've entered a statement but the only response is a `->` prompt), most likely `mysql` is waiting for the semicolon. If you don't notice what the prompt is telling you, you might sit there for a while before realizing what you need to do. Enter a semicolon to complete the statement, and `mysql` executes it:

```
mysql> SELECT USER()
    -> ;
```

```
+---------------+
| USER()        |
+---------------+
| jon@localhost |
+---------------+
```

The `'>` and `">` prompts occur during string collection (another way of saying that MySQL is waiting for completion of a string). In MySQL, you can write strings surrounded by either `'` or `"` characters (for example, `'hello'` or `"goodbye"`), and `mysql` lets you enter strings that span multiple lines. When you see a `'>` or `">` prompt, it means that you have entered a line containing a string that begins with a `'` or `"` quote character, but have not yet entered the matching quote that terminates the string. This often indicates that you have inadvertently left out a quote character. For example:

```
mysql> SELECT * FROM my_table WHERE name = 'Smith AND age < 30;
    '>
```

If you enter this `SELECT` statement, then press **Enter** and wait for the result, nothing happens. Instead of wondering why this query takes so long, notice the clue provided by the `'>` prompt. It tells you that `mysql` expects to see the rest of an unterminated string. (Do you see the error in the statement? The string `'Smith` is missing the second single quotation mark.)

At this point, what do you do? The simplest thing is to cancel the query. However, you cannot just type `\c` in this case, because `mysql` interprets it as part of the string that it is collecting. Instead, enter the closing quote character (so `mysql` knows you've finished the string), then type `\c`:

```
mysql> SELECT * FROM my_table WHERE name = 'Smith AND age < 30;
    '> '\c
mysql>
```

The prompt changes back to `mysql>`, indicating that `mysql` is ready for a new query.

The `` `> `` prompt is similar to the `'>` and `">` prompts, but indicates that you have begun but not completed a backtick-quoted identifier.

It is important to know what the `'>`, `">`, and `` `> `` prompts signify, because if you mistakenly enter an unterminated string, any further lines you type appear to be ignored by `mysql`—including a line containing `QUIT`. This can be quite confusing, especially if you do not know that you need to supply the terminating quote before you can cancel the current query.

3.3 Creating and Using a Database

Once you know how to enter SQL statements, you are ready to access a database.

Suppose that you have several pets in your home (your menagerie) and you would like to keep track of various types of information about them. You can do so by creating tables to hold your data and loading them with the desired information. Then you can answer different sorts of questions about your animals by retrieving data from the tables. This section shows you how to perform the following operations:

- Create a database
- Create a table
- Load data into the table
- Retrieve data from the table in various ways
- Use multiple tables

The menagerie database is simple (deliberately), but it is not difficult to think of real-world situations in which a similar type of database might be used. For example, a database like this could be used by a farmer to keep track of livestock, or by a veterinarian to keep track of patient records. A menagerie

distribution containing some of the queries and sample data used in the following sections can be obtained from the MySQL Web site. It is available in both compressed `tar` file and Zip formats at http://dev.mysql.com/doc/.

Use the `SHOW` statement to find out what databases currently exist on the server:

```
mysql> SHOW DATABASES;
+----------+
| Database |
+----------+
| mysql    |
| test     |
| tmp      |
+----------+
```

The `mysql` database describes user access privileges. The `test` database often is available as a workspace for users to try things out.

The list of databases displayed by the statement may be different on your machine; `SHOW DATABASES` does not show databases that you have no privileges for if you do not have the `SHOW DATABASES` privilege. See Section 13.7.5.14, “SHOW DATABASES Syntax”.

If the `test` database exists, try to access it:

```
mysql> USE test
Database changed
```

`USE`, like `QUIT`, does not require a semicolon. (You can terminate such statements with a semicolon if you like; it does no harm.) The `USE` statement is special in another way, too: it must be given on a single line.

You can use the `test` database (if you have access to it) for the examples that follow, but anything you create in that database can be removed by anyone else with access to it. For this reason, you should probably ask your MySQL administrator for permission to use a database of your own. Suppose that you want to call yours `menagerie`. The administrator needs to execute a statement like this:

```
mysql> GRANT ALL ON menagerie.* TO 'your_mysql_name'@'your_client_host';
```

where `your_mysql_name` is the MySQL user name assigned to you and `your_client_host` is the host from which you connect to the server.

3.3.1 Creating and Selecting a Database

If the administrator creates your database for you when setting up your permissions, you can begin using it. Otherwise, you need to create it yourself:

```
mysql> CREATE DATABASE menagerie;
```

Under Unix, database names are case sensitive (unlike SQL keywords), so you must always refer to your database as `menagerie`, not as `Menagerie`, `MENAGERIE`, or some other variant. This is also true for table names. (Under Windows, this restriction does not apply, although you must refer to databases and tables using the same lettercase throughout a given query. However, for a variety of reasons, the recommended best practice is always to use the same lettercase that was used when the database was created.)

Note

If you get an error such as `ERROR 1044 (42000): Access denied for user 'micah'@'localhost' to database 'menagerie'` when attempting to create a database, this means that your user account does not

have the necessary privileges to do so. Discuss this with the administrator or see Section 6.2, “The MySQL Access Privilege System”.

Creating a database does not select it for use; you must do that explicitly. To make `menagerie` the current database, use this statement:

```
mysql> USE menagerie
Database changed
```

Your database needs to be created only once, but you must select it for use each time you begin a `mysql` session. You can do this by issuing a `USE` statement as shown in the example. Alternatively, you can select the database on the command line when you invoke `mysql`. Just specify its name after any connection parameters that you might need to provide. For example:

```
shell> mysql -h host -u user -p menagerie
Enter password: ********
```

Important

`menagerie` in the command just shown is **not** your password. If you want to supply your password on the command line after the `-p` option, you must do so with no intervening space (for example, as `-pmypassword`, not as `-p mypassword`). However, putting your password on the command line is not recommended, because doing so exposes it to snooping by other users logged in on your machine.

Note

You can see at any time which database is currently selected using `SELECT DATABASE()`.

3.3.2 Creating a Table

Creating the database is the easy part, but at this point it is empty, as `SHOW TABLES` tells you:

```
mysql> SHOW TABLES;
Empty set (0.00 sec)
```

The harder part is deciding what the structure of your database should be: what tables you need and what columns should be in each of them.

You want a table that contains a record for each of your pets. This can be called the `pet` table, and it should contain, as a bare minimum, each animal's name. Because the name by itself is not very interesting, the table should contain other information. For example, if more than one person in your family keeps pets, you might want to list each animal's owner. You might also want to record some basic descriptive information such as species and sex.

How about age? That might be of interest, but it is not a good thing to store in a database. Age changes as time passes, which means you'd have to update your records often. Instead, it is better to store a fixed value such as date of birth. Then, whenever you need age, you can calculate it as the difference between the current date and the birth date. MySQL provides functions for doing date arithmetic, so this is not difficult. Storing birth date rather than age has other advantages, too:

- You can use the database for tasks such as generating reminders for upcoming pet birthdays. (If you think this type of query is somewhat silly, note that it is the same question you might ask in the context of a business database to identify clients to whom you need to send out birthday greetings in the current week or month, for that computer-assisted personal touch.)
- You can calculate age in relation to dates other than the current date. For example, if you store death date in the database, you can easily calculate how old a pet was when it died.

You can probably think of other types of information that would be useful in the `pet` table, but the ones identified so far are sufficient: name, owner, species, sex, birth, and death.

Use a `CREATE TABLE` statement to specify the layout of your table:

```
mysql> CREATE TABLE pet (name VARCHAR(20), owner VARCHAR(20),
    -> species VARCHAR(20), sex CHAR(1), birth DATE, death DATE);
```

`VARCHAR` is a good choice for the `name`, `owner`, and `species` columns because the column values vary in length. The lengths in those column definitions need not all be the same, and need not be `20`. You can normally pick any length from `1` to `65535`, whatever seems most reasonable to you. If you make a poor choice and it turns out later that you need a longer field, MySQL provides an `ALTER TABLE` statement.

Several types of values can be chosen to represent sex in animal records, such as `'m'` and `'f'`, or perhaps `'male'` and `'female'`. It is simplest to use the single characters `'m'` and `'f'`.

The use of the `DATE` data type for the `birth` and `death` columns is a fairly obvious choice.

Once you have created a table, `SHOW TABLES` should produce some output:

```
mysql> SHOW TABLES;
+---------------------+
| Tables in menagerie |
+---------------------+
| pet                 |
+---------------------+
```

To verify that your table was created the way you expected, use a `DESCRIBE` statement:

```
mysql> DESCRIBE pet;
+---------+-------------+------+-----+---------+-------+
| Field   | Type        | Null | Key | Default | Extra |
+---------+-------------+------+-----+---------+-------+
| name    | varchar(20) | YES  |     | NULL    |       |
| owner   | varchar(20) | YES  |     | NULL    |       |
| species | varchar(20) | YES  |     | NULL    |       |
| sex     | char(1)     | YES  |     | NULL    |       |
| birth   | date        | YES  |     | NULL    |       |
| death   | date        | YES  |     | NULL    |       |
+---------+-------------+------+-----+---------+-------+
```

You can use `DESCRIBE` any time, for example, if you forget the names of the columns in your table or what types they have.

For more information about MySQL data types, see Chapter 11, *Data Types*.

3.3.3 Loading Data into a Table

After creating your table, you need to populate it. The `LOAD DATA` and `INSERT` statements are useful for this.

Suppose that your pet records can be described as shown here. (Observe that MySQL expects dates in `'YYYY-MM-DD'` format; this may be different from what you are used to.)

name	owner	species	sex	birth	death
Fluffy	Harold	cat	f	1993-02-04	
Claws	Gwen	cat	m	1994-03-17	
Buffy	Harold	dog	f	1989-05-13	
Fang	Benny	dog	m	1990-08-27	
Bowser	Diane	dog	m	1979-08-31	1995-07-29

name	owner	species	sex	birth	death
Chirpy	Gwen	bird	f	1998-09-11	
Whistler	Gwen	bird		1997-12-09	
Slim	Benny	snake	m	1996-04-29	

Because you are beginning with an empty table, an easy way to populate it is to create a text file containing a row for each of your animals, then load the contents of the file into the table with a single statement.

You could create a text file `pet.txt` containing one record per line, with values separated by tabs, and given in the order in which the columns were listed in the `CREATE TABLE` statement. For missing values (such as unknown sexes or death dates for animals that are still living), you can use `NULL` values. To represent these in your text file, use `\N` (backslash, capital-N). For example, the record for Whistler the bird would look like this (where the whitespace between values is a single tab character):

```
Whistler        Gwen    bird    \N      1997-12-09      \N
```

To load the text file `pet.txt` into the `pet` table, use this statement:

```
mysql> LOAD DATA LOCAL INFILE '/path/pet.txt' INTO TABLE pet;
```

If you created the file on Windows with an editor that uses `\r\n` as a line terminator, you should use this statement instead:

```
mysql> LOAD DATA LOCAL INFILE '/path/pet.txt' INTO TABLE pet
    -> LINES TERMINATED BY '\r\n';
```

(On an Apple machine running OS X, you would likely want to use `LINES TERMINATED BY '\r'`.)

You can specify the column value separator and end of line marker explicitly in the `LOAD DATA` statement if you wish, but the defaults are tab and linefeed. These are sufficient for the statement to read the file `pet.txt` properly.

If the statement fails, it is likely that your MySQL installation does not have local file capability enabled by default. See Section 6.1.6, “Security Issues with LOAD DATA LOCAL”, for information on how to change this.

When you want to add new records one at a time, the `INSERT` statement is useful. In its simplest form, you supply values for each column, in the order in which the columns were listed in the `CREATE TABLE` statement. Suppose that Diane gets a new hamster named “Puffball.” You could add a new record using an `INSERT` statement like this:

```
mysql> INSERT INTO pet
    -> VALUES ('Puffball','Diane','hamster','f','1999-03-30',NULL);
```

String and date values are specified as quoted strings here. Also, with `INSERT`, you can insert `NULL` directly to represent a missing value. You do not use `\N` like you do with `LOAD DATA`.

From this example, you should be able to see that there would be a lot more typing involved to load your records initially using several `INSERT` statements rather than a single `LOAD DATA` statement.

3.3.4 Retrieving Information from a Table

The `SELECT` statement is used to pull information from a table. The general form of the statement is:

```
SELECT what_to_select
FROM which_table
WHERE conditions_to_satisfy;
```

what_to_select indicates what you want to see. This can be a list of columns, or * to indicate “all columns.” *which_table* indicates the table from which you want to retrieve data. The WHERE clause is optional. If it is present, *conditions_to_satisfy* specifies one or more conditions that rows must satisfy to qualify for retrieval.

3.3.4.1 Selecting All Data

The simplest form of SELECT retrieves everything from a table:

```
mysql> SELECT * FROM pet;
+----------+--------+---------+------+------------+------------+
| name     | owner  | species | sex  | birth      | death      |
+----------+--------+---------+------+------------+------------+
| Fluffy   | Harold | cat     | f    | 1993-02-04 | NULL       |
| Claws    | Gwen   | cat     | m    | 1994-03-17 | NULL       |
| Buffy    | Harold | dog     | f    | 1989-05-13 | NULL       |
| Fang     | Benny  | dog     | m    | 1990-08-27 | NULL       |
| Bowser   | Diane  | dog     | m    | 1979-08-31 | 1995-07-29 |
| Chirpy   | Gwen   | bird    | f    | 1998-09-11 | NULL       |
| Whistler | Gwen   | bird    | NULL | 1997-12-09 | NULL       |
| Slim     | Benny  | snake   | m    | 1996-04-29 | NULL       |
| Puffball | Diane  | hamster | f    | 1999-03-30 | NULL       |
+----------+--------+---------+------+------------+------------+
```

This form of SELECT is useful if you want to review your entire table, for example, after you've just loaded it with your initial data set. For example, you may happen to think that the birth date for Bowser doesn't seem quite right. Consulting your original pedigree papers, you find that the correct birth year should be 1989, not 1979.

There are at least two ways to fix this:

- Edit the file pet.txt to correct the error, then empty the table and reload it using DELETE and LOAD DATA:

  ```
  mysql> DELETE FROM pet;
  mysql> LOAD DATA LOCAL INFILE 'pet.txt' INTO TABLE pet;
  ```

  However, if you do this, you must also re-enter the record for Puffball.

- Fix only the erroneous record with an UPDATE statement:

  ```
  mysql> UPDATE pet SET birth = '1989-08-31' WHERE name = 'Bowser';
  ```

  The UPDATE changes only the record in question and does not require you to reload the table.

3.3.4.2 Selecting Particular Rows

As shown in the preceding section, it is easy to retrieve an entire table. Just omit the WHERE clause from the SELECT statement. But typically you don't want to see the entire table, particularly when it becomes large. Instead, you're usually more interested in answering a particular question, in which case you specify some constraints on the information you want. Let's look at some selection queries in terms of questions about your pets that they answer.

You can select only particular rows from your table. For example, if you want to verify the change that you made to Bowser's birth date, select Bowser's record like this:

```
mysql> SELECT * FROM pet WHERE name = 'Bowser';
+--------+-------+---------+------+------------+------------+
| name   | owner | species | sex  | birth      | death      |
+--------+-------+---------+------+------------+------------+
| Bowser | Diane | dog     | m    | 1989-08-31 | 1995-07-29 |
+--------+-------+---------+------+------------+------------+
```

The output confirms that the year is correctly recorded as 1989, not 1979.

String comparisons normally are case-insensitive, so you can specify the name as `'bowser'`, `'BOWSER'`, and so forth. The query result is the same.

You can specify conditions on any column, not just `name`. For example, if you want to know which animals were born during or after 1998, test the `birth` column:

```
mysql> SELECT * FROM pet WHERE birth >= '1998-1-1';
+----------+-------+---------+------+------------+-------+
| name     | owner | species | sex  | birth      | death |
+----------+-------+---------+------+------------+-------+
| Chirpy   | Gwen  | bird    | f    | 1998-09-11 | NULL  |
| Puffball | Diane | hamster | f    | 1999-03-30 | NULL  |
+----------+-------+---------+------+------------+-------+
```

You can combine conditions, for example, to locate female dogs:

```
mysql> SELECT * FROM pet WHERE species = 'dog' AND sex = 'f';
+-------+--------+---------+------+------------+-------+
| name  | owner  | species | sex  | birth      | death |
+-------+--------+---------+------+------------+-------+
| Buffy | Harold | dog     | f    | 1989-05-13 | NULL  |
+-------+--------+---------+------+------------+-------+
```

The preceding query uses the `AND` logical operator. There is also an `OR` operator:

```
mysql> SELECT * FROM pet WHERE species = 'snake' OR species = 'bird';
+----------+-------+---------+------+------------+-------+
| name     | owner | species | sex  | birth      | death |
+----------+-------+---------+------+------------+-------+
| Chirpy   | Gwen  | bird    | f    | 1998-09-11 | NULL  |
| Whistler | Gwen  | bird    | NULL | 1997-12-09 | NULL  |
| Slim     | Benny | snake   | m    | 1996-04-29 | NULL  |
+----------+-------+---------+------+------------+-------+
```

`AND` and `OR` may be intermixed, although `AND` has higher precedence than `OR`. If you use both operators, it is a good idea to use parentheses to indicate explicitly how conditions should be grouped:

```
mysql> SELECT * FROM pet WHERE (species = 'cat' AND sex = 'm')
    -> OR (species = 'dog' AND sex = 'f');
+-------+--------+---------+------+------------+-------+
| name  | owner  | species | sex  | birth      | death |
+-------+--------+---------+------+------------+-------+
| Claws | Gwen   | cat     | m    | 1994-03-17 | NULL  |
| Buffy | Harold | dog     | f    | 1989-05-13 | NULL  |
+-------+--------+---------+------+------------+-------+
```

3.3.4.3 Selecting Particular Columns

If you do not want to see entire rows from your table, just name the columns in which you are interested, separated by commas. For example, if you want to know when your animals were born, select the `name` and `birth` columns:

```
mysql> SELECT name, birth FROM pet;
+----------+------------+
| name     | birth      |
+----------+------------+
| Fluffy   | 1993-02-04 |
| Claws    | 1994-03-17 |
| Buffy    | 1989-05-13 |
| Fang     | 1990-08-27 |
| Bowser   | 1989-08-31 |
| Chirpy   | 1998-09-11 |
```

```
| Whistler | 1997-12-09 |
| Slim     | 1996-04-29 |
| Puffball | 1999-03-30 |
+----------+------------+
```

To find out who owns pets, use this query:

```
mysql> SELECT owner FROM pet;
+--------+
| owner  |
+--------+
| Harold |
| Gwen   |
| Harold |
| Benny  |
| Diane  |
| Gwen   |
| Gwen   |
| Benny  |
| Diane  |
+--------+
```

Notice that the query simply retrieves the `owner` column from each record, and some of them appear more than once. To minimize the output, retrieve each unique output record just once by adding the keyword `DISTINCT`:

```
mysql> SELECT DISTINCT owner FROM pet;
+--------+
| owner  |
+--------+
| Benny  |
| Diane  |
| Gwen   |
| Harold |
+--------+
```

You can use a `WHERE` clause to combine row selection with column selection. For example, to get birth dates for dogs and cats only, use this query:

```
mysql> SELECT name, species, birth FROM pet
    -> WHERE species = 'dog' OR species = 'cat';
+--------+---------+------------+
| name   | species | birth      |
+--------+---------+------------+
| Fluffy | cat     | 1993-02-04 |
| Claws  | cat     | 1994-03-17 |
| Buffy  | dog     | 1989-05-13 |
| Fang   | dog     | 1990-08-27 |
| Bowser | dog     | 1989-08-31 |
+--------+---------+------------+
```

3.3.4.4 Sorting Rows

You may have noticed in the preceding examples that the result rows are displayed in no particular order. It is often easier to examine query output when the rows are sorted in some meaningful way. To sort a result, use an `ORDER BY` clause.

Here are animal birthdays, sorted by date:

```
mysql> SELECT name, birth FROM pet ORDER BY birth;
+----------+------------+
| name     | birth      |
+----------+------------+
| Buffy    | 1989-05-13 |
| Bowser   | 1989-08-31 |
```

```
| Fang     | 1990-08-27 |
| Fluffy   | 1993-02-04 |
| Claws    | 1994-03-17 |
| Slim     | 1996-04-29 |
| Whistler | 1997-12-09 |
| Chirpy   | 1998-09-11 |
| Puffball | 1999-03-30 |
+----------+------------+
```

On character type columns, sorting—like all other comparison operations—is normally performed in a case-insensitive fashion. This means that the order is undefined for columns that are identical except for their case. You can force a case-sensitive sort for a column by using `BINARY` like so: `ORDER BY BINARY col_name`.

The default sort order is ascending, with smallest values first. To sort in reverse (descending) order, add the `DESC` keyword to the name of the column you are sorting by:

```
mysql> SELECT name, birth FROM pet ORDER BY birth DESC;
+----------+------------+
| name     | birth      |
+----------+------------+
| Puffball | 1999-03-30 |
| Chirpy   | 1998-09-11 |
| Whistler | 1997-12-09 |
| Slim     | 1996-04-29 |
| Claws    | 1994-03-17 |
| Fluffy   | 1993-02-04 |
| Fang     | 1990-08-27 |
| Bowser   | 1989-08-31 |
| Buffy    | 1989-05-13 |
+----------+------------+
```

You can sort on multiple columns, and you can sort different columns in different directions. For example, to sort by type of animal in ascending order, then by birth date within animal type in descending order (youngest animals first), use the following query:

```
mysql> SELECT name, species, birth FROM pet
    -> ORDER BY species, birth DESC;
+----------+---------+------------+
| name     | species | birth      |
+----------+---------+------------+
| Chirpy   | bird    | 1998-09-11 |
| Whistler | bird    | 1997-12-09 |
| Claws    | cat     | 1994-03-17 |
| Fluffy   | cat     | 1993-02-04 |
| Fang     | dog     | 1990-08-27 |
| Bowser   | dog     | 1989-08-31 |
| Buffy    | dog     | 1989-05-13 |
| Puffball | hamster | 1999-03-30 |
| Slim     | snake   | 1996-04-29 |
+----------+---------+------------+
```

The `DESC` keyword applies only to the column name immediately preceding it (`birth`); it does not affect the `species` column sort order.

3.3.4.5 Date Calculations

MySQL provides several functions that you can use to perform calculations on dates, for example, to calculate ages or extract parts of dates.

To determine how many years old each of your pets is, use the `TIMESTAMPDIFF()` function. Its arguments are the unit in which you want the result expressed, and the two date for which to take the difference. The following query shows, for each pet, the birth date, the current date, and the age in years. An *alias* (`age`) is used to make the final output column label more meaningful.

```
mysql> SELECT name, birth, CURDATE(),
    -> TIMESTAMPDIFF(YEAR,birth,CURDATE()) AS age
    -> FROM pet;
+----------+------------+------------+------+
| name     | birth      | CURDATE()  | age  |
+----------+------------+------------+------+
| Fluffy   | 1993-02-04 | 2003-08-19 |   10 |
| Claws    | 1994-03-17 | 2003-08-19 |    9 |
| Buffy    | 1989-05-13 | 2003-08-19 |   14 |
| Fang     | 1990-08-27 | 2003-08-19 |   12 |
| Bowser   | 1989-08-31 | 2003-08-19 |   13 |
| Chirpy   | 1998-09-11 | 2003-08-19 |    4 |
| Whistler | 1997-12-09 | 2003-08-19 |    5 |
| Slim     | 1996-04-29 | 2003-08-19 |    7 |
| Puffball | 1999-03-30 | 2003-08-19 |    4 |
+----------+------------+------------+------+
```

The query works, but the result could be scanned more easily if the rows were presented in some order. This can be done by adding an `ORDER BY name` clause to sort the output by name:

```
mysql> SELECT name, birth, CURDATE(),
    -> TIMESTAMPDIFF(YEAR,birth,CURDATE()) AS age
    -> FROM pet ORDER BY name;
+----------+------------+------------+------+
| name     | birth      | CURDATE()  | age  |
+----------+------------+------------+------+
| Bowser   | 1989-08-31 | 2003-08-19 |   13 |
| Buffy    | 1989-05-13 | 2003-08-19 |   14 |
| Chirpy   | 1998-09-11 | 2003-08-19 |    4 |
| Claws    | 1994-03-17 | 2003-08-19 |    9 |
| Fang     | 1990-08-27 | 2003-08-19 |   12 |
| Fluffy   | 1993-02-04 | 2003-08-19 |   10 |
| Puffball | 1999-03-30 | 2003-08-19 |    4 |
| Slim     | 1996-04-29 | 2003-08-19 |    7 |
| Whistler | 1997-12-09 | 2003-08-19 |    5 |
+----------+------------+------------+------+
```

To sort the output by `age` rather than `name`, just use a different `ORDER BY` clause:

```
mysql> SELECT name, birth, CURDATE(),
    -> TIMESTAMPDIFF(YEAR,birth,CURDATE()) AS age
    -> FROM pet ORDER BY age;
+----------+------------+------------+------+
| name     | birth      | CURDATE()  | age  |
+----------+------------+------------+------+
| Chirpy   | 1998-09-11 | 2003-08-19 |    4 |
| Puffball | 1999-03-30 | 2003-08-19 |    4 |
| Whistler | 1997-12-09 | 2003-08-19 |    5 |
| Slim     | 1996-04-29 | 2003-08-19 |    7 |
| Claws    | 1994-03-17 | 2003-08-19 |    9 |
| Fluffy   | 1993-02-04 | 2003-08-19 |   10 |
| Fang     | 1990-08-27 | 2003-08-19 |   12 |
| Bowser   | 1989-08-31 | 2003-08-19 |   13 |
| Buffy    | 1989-05-13 | 2003-08-19 |   14 |
+----------+------------+------------+------+
```

A similar query can be used to determine age at death for animals that have died. You determine which animals these are by checking whether the `death` value is `NULL`. Then, for those with non-`NULL` values, compute the difference between the `death` and `birth` values:

```
mysql> SELECT name, birth, death,
    -> TIMESTAMPDIFF(YEAR,birth,death) AS age
    -> FROM pet WHERE death IS NOT NULL ORDER BY age;
+--------+------------+------------+------+
| name   | birth      | death      | age  |
+--------+------------+------------+------+
| Bowser | 1989-08-31 | 1995-07-29 |    5 |
+--------+------------+------------+------+
```

The query uses `death IS NOT NULL` rather than `death <> NULL` because `NULL` is a special value that cannot be compared using the usual comparison operators. This is discussed later. See Section 3.3.4.6, “Working with NULL Values”.

What if you want to know which animals have birthdays next month? For this type of calculation, year and day are irrelevant; you simply want to extract the month part of the `birth` column. MySQL provides several functions for extracting parts of dates, such as `YEAR()`, `MONTH()`, and `DAYOFMONTH()`. `MONTH()` is the appropriate function here. To see how it works, run a simple query that displays the value of both `birth` and `MONTH(birth)`:

```
mysql> SELECT name, birth, MONTH(birth) FROM pet;
+----------+------------+--------------+
| name     | birth      | MONTH(birth) |
+----------+------------+--------------+
| Fluffy   | 1993-02-04 |            2 |
| Claws    | 1994-03-17 |            3 |
| Buffy    | 1989-05-13 |            5 |
| Fang     | 1990-08-27 |            8 |
| Bowser   | 1989-08-31 |            8 |
| Chirpy   | 1998-09-11 |            9 |
| Whistler | 1997-12-09 |           12 |
| Slim     | 1996-04-29 |            4 |
| Puffball | 1999-03-30 |            3 |
+----------+------------+--------------+
```

Finding animals with birthdays in the upcoming month is also simple. Suppose that the current month is April. Then the month value is `4` and you can look for animals born in May (month `5`) like this:

```
mysql> SELECT name, birth FROM pet WHERE MONTH(birth) = 5;
+-------+------------+
| name  | birth      |
+-------+------------+
| Buffy | 1989-05-13 |
+-------+------------+
```

There is a small complication if the current month is December. You cannot merely add one to the month number (`12`) and look for animals born in month `13`, because there is no such month. Instead, you look for animals born in January (month `1`).

You can write the query so that it works no matter what the current month is, so that you do not have to use the number for a particular month. `DATE_ADD()` enables you to add a time interval to a given date. If you add a month to the value of `CURDATE()`, then extract the month part with `MONTH()`, the result produces the month in which to look for birthdays:

```
mysql> SELECT name, birth FROM pet
    -> WHERE MONTH(birth) = MONTH(DATE_ADD(CURDATE(),INTERVAL 1 MONTH));
```

A different way to accomplish the same task is to add `1` to get the next month after the current one after using the modulo function (`MOD`) to wrap the month value to `0` if it is currently `12`:

```
mysql> SELECT name, birth FROM pet
    -> WHERE MONTH(birth) = MOD(MONTH(CURDATE()), 12) + 1;
```

`MONTH()` returns a number between `1` and `12`. And `MOD(something,12)` returns a number between `0` and `11`. So the addition has to be after the `MOD()`, otherwise we would go from November (`11`) to January (`1`).

3.3.4.6 Working with NULL Values

The `NULL` value can be surprising until you get used to it. Conceptually, `NULL` means “a missing unknown value” and it is treated somewhat differently from other values.

To test for NULL, use the IS NULL and IS NOT NULL operators, as shown here:

```
mysql> SELECT 1 IS NULL, 1 IS NOT NULL;
+-----------+---------------+
| 1 IS NULL | 1 IS NOT NULL |
+-----------+---------------+
|         0 |             1 |
+-----------+---------------+
```

You cannot use arithmetic comparison operators such as =, <, or <> to test for NULL. To demonstrate this for yourself, try the following query:

```
mysql> SELECT 1 = NULL, 1 <> NULL, 1 < NULL, 1 > NULL;
+----------+-----------+----------+----------+
| 1 = NULL | 1 <> NULL | 1 < NULL | 1 > NULL |
+----------+-----------+----------+----------+
|     NULL |      NULL |     NULL |     NULL |
+----------+-----------+----------+----------+
```

Because the result of any arithmetic comparison with NULL is also NULL, you cannot obtain any meaningful results from such comparisons.

In MySQL, 0 or NULL means false and anything else means true. The default truth value from a boolean operation is 1.

This special treatment of NULL is why, in the previous section, it was necessary to determine which animals are no longer alive using death IS NOT NULL instead of death <> NULL.

Two NULL values are regarded as equal in a GROUP BY.

When doing an ORDER BY, NULL values are presented first if you do ORDER BY ... ASC and last if you do ORDER BY ... DESC.

A common error when working with NULL is to assume that it is not possible to insert a zero or an empty string into a column defined as NOT NULL, but this is not the case. These are in fact values, whereas NULL means “not having a value.” You can test this easily enough by using IS [NOT] NULL as shown:

```
mysql> SELECT 0 IS NULL, 0 IS NOT NULL, '' IS NULL, '' IS NOT NULL;
+-----------+---------------+------------+----------------+
| 0 IS NULL | 0 IS NOT NULL | '' IS NULL | '' IS NOT NULL |
+-----------+---------------+------------+----------------+
|         0 |             1 |          0 |              1 |
+-----------+---------------+------------+----------------+
```

Thus it is entirely possible to insert a zero or empty string into a NOT NULL column, as these are in fact NOT NULL. See Section B.5.4.3, “Problems with NULL Values”.

3.3.4.7 Pattern Matching

MySQL provides standard SQL pattern matching as well as a form of pattern matching based on extended regular expressions similar to those used by Unix utilities such as vi, grep, and sed.

SQL pattern matching enables you to use _ to match any single character and % to match an arbitrary number of characters (including zero characters). In MySQL, SQL patterns are case-insensitive by default. Some examples are shown here. You do not use = or <> when you use SQL patterns; use the LIKE or NOT LIKE comparison operators instead.

To find names beginning with b:

```
mysql> SELECT * FROM pet WHERE name LIKE 'b%';
+--------+--------+---------+------+------------+------------+
```

```
| name   | owner  | species | sex  | birth      | death      |
+--------+--------+---------+------+------------+------------+
| Buffy  | Harold | dog     | f    | 1989-05-13 | NULL       |
| Bowser | Diane  | dog     | m    | 1989-08-31 | 1995-07-29 |
+--------+--------+---------+------+------------+------------+
```

To find names ending with `fy`:

```
mysql> SELECT * FROM pet WHERE name LIKE '%fy';
+--------+--------+---------+------+------------+-------+
| name   | owner  | species | sex  | birth      | death |
+--------+--------+---------+------+------------+-------+
| Fluffy | Harold | cat     | f    | 1993-02-04 | NULL  |
| Buffy  | Harold | dog     | f    | 1989-05-13 | NULL  |
+--------+--------+---------+------+------------+-------+
```

To find names containing a `w`:

```
mysql> SELECT * FROM pet WHERE name LIKE '%w%';
+----------+-------+---------+------+------------+------------+
| name     | owner | species | sex  | birth      | death      |
+----------+-------+---------+------+------------+------------+
| Claws    | Gwen  | cat     | m    | 1994-03-17 | NULL       |
| Bowser   | Diane | dog     | m    | 1989-08-31 | 1995-07-29 |
| Whistler | Gwen  | bird    | NULL | 1997-12-09 | NULL       |
+----------+-------+---------+------+------------+------------+
```

To find names containing exactly five characters, use five instances of the _ pattern character:

```
mysql> SELECT * FROM pet WHERE name LIKE '_____';
+-------+--------+---------+------+------------+-------+
| name  | owner  | species | sex  | birth      | death |
+-------+--------+---------+------+------------+-------+
| Claws | Gwen   | cat     | m    | 1994-03-17 | NULL  |
| Buffy | Harold | dog     | f    | 1989-05-13 | NULL  |
+-------+--------+---------+------+------------+-------+
```

The other type of pattern matching provided by MySQL uses extended regular expressions. When you test for a match for this type of pattern, use the `REGEXP` and `NOT REGEXP` operators (or `RLIKE` and `NOT RLIKE`, which are synonyms).

The following list describes some characteristics of extended regular expressions:

- `.` matches any single character.
- A character class `[...]` matches any character within the brackets. For example, `[abc]` matches `a`, `b`, or `c`. To name a range of characters, use a dash. `[a-z]` matches any letter, whereas `[0-9]` matches any digit.
- `*` matches zero or more instances of the thing preceding it. For example, `x*` matches any number of `x` characters, `[0-9]*` matches any number of digits, and `.*` matches any number of anything.
- A `REGEXP` pattern match succeeds if the pattern matches anywhere in the value being tested. (This differs from a `LIKE` pattern match, which succeeds only if the pattern matches the entire value.)
- To anchor a pattern so that it must match the beginning or end of the value being tested, use `^` at the beginning or `$` at the end of the pattern.

To demonstrate how extended regular expressions work, the `LIKE` queries shown previously are rewritten here to use `REGEXP`.

To find names beginning with `b`, use `^` to match the beginning of the name:

```
mysql> SELECT * FROM pet WHERE name REGEXP '^b';
+--------+--------+---------+------+------------+------------+
| name   | owner  | species | sex  | birth      | death      |
+--------+--------+---------+------+------------+------------+
| Buffy  | Harold | dog     | f    | 1989-05-13 | NULL       |
| Bowser | Diane  | dog     | m    | 1989-08-31 | 1995-07-29 |
+--------+--------+---------+------+------------+------------+
```

If you really want to force a `REGEXP` comparison to be case sensitive, use the `BINARY` keyword to make one of the strings a binary string. This query matches only lowercase `b` at the beginning of a name:

```
mysql> SELECT * FROM pet WHERE name REGEXP BINARY '^b';
```

To find names ending with `fy`, use `$` to match the end of the name:

```
mysql> SELECT * FROM pet WHERE name REGEXP 'fy$';
+--------+--------+---------+------+------------+-------+
| name   | owner  | species | sex  | birth      | death |
+--------+--------+---------+------+------------+-------+
| Fluffy | Harold | cat     | f    | 1993-02-04 | NULL  |
| Buffy  | Harold | dog     | f    | 1989-05-13 | NULL  |
+--------+--------+---------+------+------------+-------+
```

To find names containing a `w`, use this query:

```
mysql> SELECT * FROM pet WHERE name REGEXP 'w';
+----------+-------+---------+------+------------+------------+
| name     | owner | species | sex  | birth      | death      |
+----------+-------+---------+------+------------+------------+
| Claws    | Gwen  | cat     | m    | 1994-03-17 | NULL       |
| Bowser   | Diane | dog     | m    | 1989-08-31 | 1995-07-29 |
| Whistler | Gwen  | bird    | NULL | 1997-12-09 | NULL       |
+----------+-------+---------+------+------------+------------+
```

Because a regular expression pattern matches if it occurs anywhere in the value, it is not necessary in the previous query to put a wildcard on either side of the pattern to get it to match the entire value like it would be if you used an SQL pattern.

To find names containing exactly five characters, use `^` and `$` to match the beginning and end of the name, and five instances of `.` in between:

```
mysql> SELECT * FROM pet WHERE name REGEXP '^.....$';
+-------+--------+---------+------+------------+-------+
| name  | owner  | species | sex  | birth      | death |
+-------+--------+---------+------+------------+-------+
| Claws | Gwen   | cat     | m    | 1994-03-17 | NULL  |
| Buffy | Harold | dog     | f    | 1989-05-13 | NULL  |
+-------+--------+---------+------+------------+-------+
```

You could also write the previous query using the `{n}` (“repeat-*n*-times”) operator:

```
mysql> SELECT * FROM pet WHERE name REGEXP '^.{5}$';
+-------+--------+---------+------+------------+-------+
| name  | owner  | species | sex  | birth      | death |
+-------+--------+---------+------+------------+-------+
| Claws | Gwen   | cat     | m    | 1994-03-17 | NULL  |
| Buffy | Harold | dog     | f    | 1989-05-13 | NULL  |
+-------+--------+---------+------+------------+-------+
```

Section 12.5.2, “Regular Expressions”, provides more information about the syntax for regular expressions.

3.3.4.8 Counting Rows

Databases are often used to answer the question, “How often does a certain type of data occur in a table?” For example, you might want to know how many pets you have, or how many pets each owner has, or you might want to perform various kinds of census operations on your animals.

Counting the total number of animals you have is the same question as “How many rows are in the `pet` table?” because there is one record per pet. `COUNT(*)` counts the number of rows, so the query to count your animals looks like this:

```
mysql> SELECT COUNT(*) FROM pet;
+----------+
| COUNT(*) |
+----------+
|        9 |
+----------+
```

Earlier, you retrieved the names of the people who owned pets. You can use `COUNT()` if you want to find out how many pets each owner has:

```
mysql> SELECT owner, COUNT(*) FROM pet GROUP BY owner;
+--------+----------+
| owner  | COUNT(*) |
+--------+----------+
| Benny  |        2 |
| Diane  |        2 |
| Gwen   |        3 |
| Harold |        2 |
+--------+----------+
```

The preceding query uses `GROUP BY` to group all records for each `owner`. The use of `COUNT()` in conjunction with `GROUP BY` is useful for characterizing your data under various groupings. The following examples show different ways to perform animal census operations.

Number of animals per species:

```
mysql> SELECT species, COUNT(*) FROM pet GROUP BY species;
+---------+----------+
| species | COUNT(*) |
+---------+----------+
| bird    |        2 |
| cat     |        2 |
| dog     |        3 |
| hamster |        1 |
| snake   |        1 |
+---------+----------+
```

Number of animals per sex:

```
mysql> SELECT sex, COUNT(*) FROM pet GROUP BY sex;
+------+----------+
| sex  | COUNT(*) |
+------+----------+
| NULL |        1 |
| f    |        4 |
| m    |        4 |
+------+----------+
```

(In this output, `NULL` indicates that the sex is unknown.)

Number of animals per combination of species and sex:

```
mysql> SELECT species, sex, COUNT(*) FROM pet GROUP BY species, sex;
+---------+------+----------+
| species | sex  | COUNT(*) |
+---------+------+----------+
```

```
| bird    | NULL |        1 |
| bird    | f    |        1 |
| cat     | f    |        1 |
| cat     | m    |        1 |
| dog     | f    |        1 |
| dog     | m    |        2 |
| hamster | f    |        1 |
| snake   | m    |        1 |
+---------+------+----------+
```

You need not retrieve an entire table when you use `COUNT()`. For example, the previous query, when performed just on dogs and cats, looks like this:

```
mysql> SELECT species, sex, COUNT(*) FROM pet
    -> WHERE species = 'dog' OR species = 'cat'
    -> GROUP BY species, sex;
+---------+------+----------+
| species | sex  | COUNT(*) |
+---------+------+----------+
| cat     | f    |        1 |
| cat     | m    |        1 |
| dog     | f    |        1 |
| dog     | m    |        2 |
+---------+------+----------+
```

Or, if you wanted the number of animals per sex only for animals whose sex is known:

```
mysql> SELECT species, sex, COUNT(*) FROM pet
    -> WHERE sex IS NOT NULL
    -> GROUP BY species, sex;
+---------+------+----------+
| species | sex  | COUNT(*) |
+---------+------+----------+
| bird    | f    |        1 |
| cat     | f    |        1 |
| cat     | m    |        1 |
| dog     | f    |        1 |
| dog     | m    |        2 |
| hamster | f    |        1 |
| snake   | m    |        1 |
+---------+------+----------+
```

If you name columns to select in addition to the `COUNT()` value, a `GROUP BY` clause should be present that names those same columns. Otherwise, the following occurs:

- If the `ONLY_FULL_GROUP_BY` SQL mode is enabled, an error occurs:

  ```
  mysql> SET sql_mode = 'ONLY_FULL_GROUP_BY';
  Query OK, 0 rows affected (0.00 sec)

  mysql> SELECT owner, COUNT(*) FROM pet;
  ERROR 1140 (42000): In aggregated query without GROUP BY, expression
  #1 of SELECT list contains nonaggregated column 'menagerie.pet.owner';
  this is incompatible with sql_mode=only_full_group_by
  ```

- If `ONLY_FULL_GROUP_BY` is not enabled, the query is processed by treating all rows as a single group, but the value selected for each named column is indeterminate. The server is free to select the value from any row:

  ```
  mysql> SET sql_mode = '';
  Query OK, 0 rows affected (0.00 sec)

  mysql> SELECT owner, COUNT(*) FROM pet;
  +--------+----------+
  | owner  | COUNT(*) |
  +--------+----------+
  ```

```
| Harold |        8 |
+--------+----------+
1 row in set (0.00 sec)
```

See also Section 12.18.3, “MySQL Handling of GROUP BY”. See Section 12.18.1, “Aggregate (GROUP BY) Function Descriptions” for information about `COUNT(expr)` behavior and related optimizations.

3.3.4.9 Using More Than one Table

The `pet` table keeps track of which pets you have. If you want to record other information about them, such as events in their lives like visits to the vet or when litters are born, you need another table. What should this table look like? It needs to contain the following information:

- The pet name so that you know which animal each event pertains to.
- A date so that you know when the event occurred.
- A field to describe the event.
- An event type field, if you want to be able to categorize events.

Given these considerations, the `CREATE TABLE` statement for the `event` table might look like this:

```
mysql> CREATE TABLE event (name VARCHAR(20), date DATE,
    -> type VARCHAR(15), remark VARCHAR(255));
```

As with the `pet` table, it is easiest to load the initial records by creating a tab-delimited text file containing the following information.

name	date	type	remark
Fluffy	1995-05-15	litter	4 kittens, 3 female, 1 male
Buffy	1993-06-23	litter	5 puppies, 2 female, 3 male
Buffy	1994-06-19	litter	3 puppies, 3 female
Chirpy	1999-03-21	vet	needed beak straightened
Slim	1997-08-03	vet	broken rib
Bowser	1991-10-12	kennel	
Fang	1991-10-12	kennel	
Fang	1998-08-28	birthday	Gave him a new chew toy
Claws	1998-03-17	birthday	Gave him a new flea collar
Whistler	1998-12-09	birthday	First birthday

Load the records like this:

```
mysql> LOAD DATA LOCAL INFILE 'event.txt' INTO TABLE event;
```

Based on what you have learned from the queries that you have run on the `pet` table, you should be able to perform retrievals on the records in the `event` table; the principles are the same. But when is the `event` table by itself insufficient to answer questions you might ask?

Suppose that you want to find out the ages at which each pet had its litters. We saw earlier how to calculate ages from two dates. The litter date of the mother is in the `event` table, but to calculate her age on that date you need her birth date, which is stored in the `pet` table. This means the query requires both tables:

```
mysql> SELECT pet.name,
```

```
    -> TIMESTAMPDIFF(YEAR,birth,date) AS age,
    -> remark
    -> FROM pet INNER JOIN event
    ->   ON pet.name = event.name
    -> WHERE event.type = 'litter';
+--------+------+-----------------------------+
| name   | age  | remark                      |
+--------+------+-----------------------------+
| Fluffy |    2 | 4 kittens, 3 female, 1 male |
| Buffy  |    4 | 5 puppies, 2 female, 3 male |
| Buffy  |    5 | 3 puppies, 3 female         |
+--------+------+-----------------------------+
```

There are several things to note about this query:

- The `FROM` clause joins two tables because the query needs to pull information from both of them.
- When combining (joining) information from multiple tables, you need to specify how records in one table can be matched to records in the other. This is easy because they both have a `name` column. The query uses an `ON` clause to match up records in the two tables based on the `name` values.

  The query uses an `INNER JOIN` to combine the tables. An `INNER JOIN` permits rows from either table to appear in the result if and only if both tables meet the conditions specified in the `ON` clause. In this example, the `ON` clause specifies that the `name` column in the `pet` table must match the `name` column in the `event` table. If a name appears in one table but not the other, the row will not appear in the result because the condition in the `ON` clause fails.
- Because the `name` column occurs in both tables, you must be specific about which table you mean when referring to the column. This is done by prepending the table name to the column name.

You need not have two different tables to perform a join. Sometimes it is useful to join a table to itself, if you want to compare records in a table to other records in that same table. For example, to find breeding pairs among your pets, you can join the `pet` table with itself to produce candidate pairs of males and females of like species:

```
mysql> SELECT p1.name, p1.sex, p2.name, p2.sex, p1.species
    -> FROM pet AS p1 INNER JOIN pet AS p2
    ->   ON p1.species = p2.species AND p1.sex = 'f' AND p2.sex = 'm';
+--------+------+--------+------+---------+
| name   | sex  | name   | sex  | species |
+--------+------+--------+------+---------+
| Fluffy | f    | Claws  | m    | cat     |
| Buffy  | f    | Fang   | m    | dog     |
| Buffy  | f    | Bowser | m    | dog     |
+--------+------+--------+------+---------+
```

In this query, we specify aliases for the table name to refer to the columns and keep straight which instance of the table each column reference is associated with.

3.4 Getting Information About Databases and Tables

What if you forget the name of a database or table, or what the structure of a given table is (for example, what its columns are called)? MySQL addresses this problem through several statements that provide information about the databases and tables it supports.

You have previously seen `SHOW DATABASES`, which lists the databases managed by the server. To find out which database is currently selected, use the `DATABASE()` function:

```
mysql> SELECT DATABASE();
+------------+
| DATABASE() |
+------------+
| menagerie  |
+------------+
```

If you have not yet selected any database, the result is `NULL`.

To find out what tables the default database contains (for example, when you are not sure about the name of a table), use this statement:

```
mysql> SHOW TABLES;
+--------------------+
| Tables_in_menagerie |
+--------------------+
| event              |
| pet                |
+--------------------+
```

The name of the column in the output produced by this statement is always `Tables_in_db_name`, where *db_name* is the name of the database. See Section 13.7.5.37, “SHOW TABLES Syntax”, for more information.

If you want to find out about the structure of a table, the `DESCRIBE` statement is useful; it displays information about each of a table's columns:

```
mysql> DESCRIBE pet;
+---------+-------------+------+-----+---------+-------+
| Field   | Type        | Null | Key | Default | Extra |
+---------+-------------+------+-----+---------+-------+
| name    | varchar(20) | YES  |     | NULL    |       |
| owner   | varchar(20) | YES  |     | NULL    |       |
| species | varchar(20) | YES  |     | NULL    |       |
| sex     | char(1)     | YES  |     | NULL    |       |
| birth   | date        | YES  |     | NULL    |       |
| death   | date        | YES  |     | NULL    |       |
+---------+-------------+------+-----+---------+-------+
```

`Field` indicates the column name, `Type` is the data type for the column, `NULL` indicates whether the column can contain `NULL` values, `Key` indicates whether the column is indexed, and `Default` specifies the column's default value. `Extra` displays special information about columns: If a column was created with the `AUTO_INCREMENT` option, the value will be `auto_increment` rather than empty.

`DESC` is a short form of `DESCRIBE`. See Section 13.8.1, “DESCRIBE Syntax”, for more information.

You can obtain the `CREATE TABLE` statement necessary to create an existing table using the `SHOW CREATE TABLE` statement. See Section 13.7.5.10, “SHOW CREATE TABLE Syntax”.

If you have indexes on a table, `SHOW INDEX FROM tbl_name` produces information about them. See Section 13.7.5.22, “SHOW INDEX Syntax”, for more about this statement.

3.5 Using mysql in Batch Mode

In the previous sections, you used `mysql` interactively to enter statements and view the results. You can also run `mysql` in batch mode. To do this, put the statements you want to run in a file, then tell `mysql` to read its input from the file:

```
shell> mysql < batch-file
```

If you are running `mysql` under Windows and have some special characters in the file that cause problems, you can do this:

```
C:\> mysql -e "source batch-file"
```

If you need to specify connection parameters on the command line, the command might look like this:

```
shell> mysql -h host -u user -p < batch-file
Enter password: ********
```

When you use `mysql` this way, you are creating a script file, then executing the script.

If you want the script to continue even if some of the statements in it produce errors, you should use the `--force` command-line option.

Why use a script? Here are a few reasons:

- If you run a query repeatedly (say, every day or every week), making it a script enables you to avoid retyping it each time you execute it.
- You can generate new queries from existing ones that are similar by copying and editing script files.
- Batch mode can also be useful while you're developing a query, particularly for multiple-line statements or multiple-statement sequences. If you make a mistake, you don't have to retype everything. Just edit your script to correct the error, then tell `mysql` to execute it again.
- If you have a query that produces a lot of output, you can run the output through a pager rather than watching it scroll off the top of your screen:

```
shell> mysql < batch-file | more
```

- You can catch the output in a file for further processing:

```
shell> mysql < batch-file > mysql.out
```

- You can distribute your script to other people so that they can also run the statements.
- Some situations do not allow for interactive use, for example, when you run a query from a `cron` job. In this case, you must use batch mode.

The default output format is different (more concise) when you run `mysql` in batch mode than when you use it interactively. For example, the output of `SELECT DISTINCT species FROM pet` looks like this when `mysql` is run interactively:

```
+---------+
| species |
+---------+
| bird    |
| cat     |
| dog     |
| hamster |
| snake   |
+---------+
```

In batch mode, the output looks like this instead:

```
species
bird
cat
dog
hamster
snake
```

If you want to get the interactive output format in batch mode, use `mysql -t`. To echo to the output the statements that are executed, use `mysql -v`.

You can also use scripts from the `mysql` prompt by using the `source` command or `\.` command:

```
mysql> source filename;
mysql> \. filename
```

See Section 4.5.1.5, “Executing SQL Statements from a Text File”, for more information.

3.6 Examples of Common Queries

Here are examples of how to solve some common problems with MySQL.

Some of the examples use the table `shop` to hold the price of each article (item number) for certain traders (dealers). Supposing that each trader has a single fixed price per article, then (`article`, `dealer`) is a primary key for the records.

Start the command-line tool `mysql` and select a database:

```
shell> mysql your-database-name
```

(In most MySQL installations, you can use the database named `test`).

You can create and populate the example table with these statements:

```
CREATE TABLE shop (
    article INT(4) UNSIGNED ZEROFILL DEFAULT '0000' NOT NULL,
    dealer  CHAR(20)                 DEFAULT ''     NOT NULL,
    price   DOUBLE(16,2)             DEFAULT '0.00' NOT NULL,
    PRIMARY KEY(article, dealer));
INSERT INTO shop VALUES
    (1,'A',3.45),(1,'B',3.99),(2,'A',10.99),(3,'B',1.45),
    (3,'C',1.69),(3,'D',1.25),(4,'D',19.95);
```

After issuing the statements, the table should have the following contents:

```
SELECT * FROM shop;

+---------+--------+-------+
| article | dealer | price |
+---------+--------+-------+
|    0001 | A      |  3.45 |
|    0001 | B      |  3.99 |
|    0002 | A      | 10.99 |
|    0003 | B      |  1.45 |
|    0003 | C      |  1.69 |
|    0003 | D      |  1.25 |
|    0004 | D      | 19.95 |
+---------+--------+-------+
```

3.6.1 The Maximum Value for a Column

“What is the highest item number?”

```
SELECT MAX(article) AS article FROM shop;

+---------+
| article |
+---------+
|       4 |
+---------+
```

3.6.2 The Row Holding the Maximum of a Certain Column

Task: Find the number, dealer, and price of the most expensive article.

This is easily done with a subquery:

```
SELECT article, dealer, price
FROM   shop
WHERE  price=(SELECT MAX(price) FROM shop);
```

```
+---------+--------+-------+
| article | dealer | price |
+---------+--------+-------+
|    0004 | D      | 19.95 |
+---------+--------+-------+
```

Other solutions are to use a `LEFT JOIN` or to sort all rows descending by price and get only the first row using the MySQL-specific `LIMIT` clause:

```
SELECT s1.article, s1.dealer, s1.price
FROM shop s1
LEFT JOIN shop s2 ON s1.price < s2.price
WHERE s2.article IS NULL;

SELECT article, dealer, price
FROM shop
ORDER BY price DESC
LIMIT 1;
```

Note

If there were several most expensive articles, each with a price of 19.95, the `LIMIT` solution would show only one of them.

3.6.3 Maximum of Column per Group

Task: Find the highest price per article.

```
SELECT article, MAX(price) AS price
FROM   shop
GROUP BY article;

+---------+-------+
| article | price |
+---------+-------+
|    0001 |  3.99 |
|    0002 | 10.99 |
|    0003 |  1.69 |
|    0004 | 19.95 |
+---------+-------+
```

3.6.4 The Rows Holding the Group-wise Maximum of a Certain Column

Task: For each article, find the dealer or dealers with the most expensive price.

This problem can be solved with a subquery like this one:

```
SELECT article, dealer, price
FROM   shop s1
WHERE  price=(SELECT MAX(s2.price)
              FROM shop s2
              WHERE s1.article = s2.article);

+---------+--------+-------+
| article | dealer | price |
+---------+--------+-------+
|    0001 | B      |  3.99 |
|    0002 | A      | 10.99 |
|    0003 | C      |  1.69 |
|    0004 | D      | 19.95 |
+---------+--------+-------+
```

The preceding example uses a correlated subquery, which can be inefficient (see Section 13.2.11.7, “Correlated Subqueries”). Other possibilities for solving the problem are to use an uncorrelated subquery in the `FROM` clause or a `LEFT JOIN`.

Uncorrelated subquery:

```
SELECT s1.article, dealer, s1.price
FROM shop s1
JOIN (
  SELECT article, MAX(price) AS price
  FROM shop
  GROUP BY article) AS s2
  ON s1.article = s2.article AND s1.price = s2.price;
```

`LEFT JOIN`:

```
SELECT s1.article, s1.dealer, s1.price
FROM shop s1
LEFT JOIN shop s2 ON s1.article = s2.article AND s1.price < s2.price
WHERE s2.article IS NULL;
```

The `LEFT JOIN` works on the basis that when `s1.price` is at its maximum value, there is no `s2.price` with a greater value and the `s2` rows values will be `NULL`. See Section 13.2.10.2, “JOIN Syntax”.

3.6.5 Using User-Defined Variables

You can employ MySQL user variables to remember results without having to store them in temporary variables in the client. (See Section 9.4, “User-Defined Variables”.)

For example, to find the articles with the highest and lowest price you can do this:

```
mysql> SELECT @min_price:=MIN(price),@max_price:=MAX(price) FROM shop;
mysql> SELECT * FROM shop WHERE price=@min_price OR price=@max_price;
+---------+--------+-------+
| article | dealer | price |
+---------+--------+-------+
|    0003 | D      |  1.25 |
|    0004 | D      | 19.95 |
+---------+--------+-------+
```

Note

It is also possible to store the name of a database object such as a table or a column in a user variable and then to use this variable in an SQL statement; however, this requires the use of a prepared statement. See Section 13.5, “Prepared SQL Statement Syntax”, for more information.

3.6.6 Using Foreign Keys

In MySQL, `InnoDB` tables support checking of foreign key constraints. See Chapter 15, *The InnoDB Storage Engine*, and Section 1.8.2.3, “Foreign Key Differences”.

A foreign key constraint is not required merely to join two tables. For storage engines other than `InnoDB`, it is possible when defining a column to use a `REFERENCES` *`tbl_name(col_name)`* clause, which has no actual effect, and *serves only as a memo or comment to you that the column which you are currently defining is intended to refer to a column in another table*. It is extremely important to realize when using this syntax that:

- MySQL does not perform any sort of `CHECK` to make sure that *`col_name`* actually exists in *`tbl_name`* (or even that *`tbl_name`* itself exists).
- MySQL does not perform any sort of action on *`tbl_name`* such as deleting rows in response to actions taken on rows in the table which you are defining; in other words, this syntax induces no `ON DELETE` or `ON UPDATE` behavior whatsoever. (Although you can write an `ON DELETE` or `ON UPDATE` clause as part of the `REFERENCES` clause, it is also ignored.)

- This syntax creates a *column*; it does **not** create any sort of index or key.

You can use a column so created as a join column, as shown here:

```
CREATE TABLE person (
    id SMALLINT UNSIGNED NOT NULL AUTO_INCREMENT,
    name CHAR(60) NOT NULL,
    PRIMARY KEY (id)
);

CREATE TABLE shirt (
    id SMALLINT UNSIGNED NOT NULL AUTO_INCREMENT,
    style ENUM('t-shirt', 'polo', 'dress') NOT NULL,
    color ENUM('red', 'blue', 'orange', 'white', 'black') NOT NULL,
    owner SMALLINT UNSIGNED NOT NULL REFERENCES person(id),
    PRIMARY KEY (id)
);

INSERT INTO person VALUES (NULL, 'Antonio Paz');

SELECT @last := LAST_INSERT_ID();

INSERT INTO shirt VALUES
(NULL, 'polo', 'blue', @last),
(NULL, 'dress', 'white', @last),
(NULL, 't-shirt', 'blue', @last);

INSERT INTO person VALUES (NULL, 'Lilliana Angelovska');

SELECT @last := LAST_INSERT_ID();

INSERT INTO shirt VALUES
(NULL, 'dress', 'orange', @last),
(NULL, 'polo', 'red', @last),
(NULL, 'dress', 'blue', @last),
(NULL, 't-shirt', 'white', @last);

SELECT * FROM person;
+----+--------------------+
| id | name               |
+----+--------------------+
|  1 | Antonio Paz        |
|  2 | Lilliana Angelovska |
+----+--------------------+

SELECT * FROM shirt;
+----+---------+--------+-------+
| id | style   | color  | owner |
+----+---------+--------+-------+
|  1 | polo    | blue   |     1 |
|  2 | dress   | white  |     1 |
|  3 | t-shirt | blue   |     1 |
|  4 | dress   | orange |     2 |
|  5 | polo    | red    |     2 |
|  6 | dress   | blue   |     2 |
|  7 | t-shirt | white  |     2 |
+----+---------+--------+-------+


SELECT s.* FROM person p INNER JOIN shirt s
   ON s.owner = p.id
 WHERE p.name LIKE 'Lilliana%'
   AND s.color <> 'white';

+----+-------+--------+-------+
| id | style | color  | owner |
+----+-------+--------+-------+
|  4 | dress | orange |     2 |
|  5 | polo  | red    |     2 |
|  6 | dress | blue   |     2 |
+----+-------+--------+-------+
```

When used in this fashion, the `REFERENCES` clause is not displayed in the output of `SHOW CREATE TABLE` or `DESCRIBE`:

```
SHOW CREATE TABLE shirt\G
*************************** 1. row ***************************
Table: shirt
Create Table: CREATE TABLE `shirt` (
`id` smallint(5) unsigned NOT NULL auto_increment,
`style` enum('t-shirt','polo','dress') NOT NULL,
`color` enum('red','blue','orange','white','black') NOT NULL,
`owner` smallint(5) unsigned NOT NULL,
PRIMARY KEY  (`id`)
) ENGINE=MyISAM DEFAULT CHARSET=utf8mb4
```

The use of `REFERENCES` in this way as a comment or “reminder” in a column definition works with `MyISAM` tables.

3.6.7 Searching on Two Keys

An `OR` using a single key is well optimized, as is the handling of `AND`.

The one tricky case is that of searching on two different keys combined with `OR`:

```
SELECT field1_index, field2_index FROM test_table
WHERE field1_index = '1' OR  field2_index = '1'
```

This case is optimized. See Section 8.2.1.3, “Index Merge Optimization”.

You can also solve the problem efficiently by using a `UNION` that combines the output of two separate `SELECT` statements. See Section 13.2.10.3, “UNION Syntax”.

Each `SELECT` searches only one key and can be optimized:

```
SELECT field1_index, field2_index
    FROM test_table WHERE field1_index = '1'
UNION
SELECT field1_index, field2_index
    FROM test_table WHERE field2_index = '1';
```

3.6.8 Calculating Visits Per Day

The following example shows how you can use the bit group functions to calculate the number of days per month a user has visited a Web page.

```
CREATE TABLE t1 (year YEAR(4), month INT(2) UNSIGNED ZEROFILL,
             day INT(2) UNSIGNED ZEROFILL);
INSERT INTO t1 VALUES(2000,1,1),(2000,1,20),(2000,1,30),(2000,2,2),
            (2000,2,23),(2000,2,23);
```

The example table contains year-month-day values representing visits by users to the page. To determine how many different days in each month these visits occur, use this query:

```
SELECT year,month,BIT_COUNT(BIT_OR(1<<day)) AS days FROM t1
       GROUP BY year,month;
```

Which returns:

```
+------+-------+------+
| year | month | days |
+------+-------+------+
| 2000 |    01 |    3 |
| 2000 |    02 |    2 |
```

```
+------+-------+------+
```

The query calculates how many different days appear in the table for each year/month combination, with automatic removal of duplicate entries.

3.6.9 Using AUTO_INCREMENT

The `AUTO_INCREMENT` attribute can be used to generate a unique identity for new rows:

```
CREATE TABLE animals (
     id MEDIUMINT NOT NULL AUTO_INCREMENT,
     name CHAR(30) NOT NULL,
     PRIMARY KEY (id)
);

INSERT INTO animals (name) VALUES
    ('dog'),('cat'),('penguin'),
    ('lax'),('whale'),('ostrich');

SELECT * FROM animals;
```

Which returns:

```
+----+---------+
| id | name    |
+----+---------+
|  1 | dog     |
|  2 | cat     |
|  3 | penguin |
|  4 | lax     |
|  5 | whale   |
|  6 | ostrich |
+----+---------+
```

No value was specified for the `AUTO_INCREMENT` column, so MySQL assigned sequence numbers automatically. You can also explicitly assign 0 to the column to generate sequence numbers, unless the `NO_AUTO_VALUE_ON_ZERO` SQL mode is enabled. If the column is declared `NOT NULL`, it is also possible to assign `NULL` to the column to generate sequence numbers. When you insert any other value into an `AUTO_INCREMENT` column, the column is set to that value and the sequence is reset so that the next automatically generated value follows sequentially from the largest column value.

You can retrieve the most recent automatically generated `AUTO_INCREMENT` value with the `LAST_INSERT_ID()` SQL function or the `mysql_insert_id()` C API function. These functions are connection-specific, so their return values are not affected by another connection which is also performing inserts.

Use the smallest integer data type for the `AUTO_INCREMENT` column that is large enough to hold the maximum sequence value you will need. When the column reaches the upper limit of the data type, the next attempt to generate a sequence number fails. Use the `UNSIGNED` attribute if possible to allow a greater range. For example, if you use `TINYINT`, the maximum permissible sequence number is 127. For `TINYINT UNSIGNED`, the maximum is 255. See Section 11.2.1, “Integer Types (Exact Value) - INTEGER, INT, SMALLINT, TINYINT, MEDIUMINT, BIGINT” for the ranges of all the integer types.

Note

For a multiple-row insert, `LAST_INSERT_ID()` and `mysql_insert_id()` actually return the `AUTO_INCREMENT` key from the *first* of the inserted rows. This enables multiple-row inserts to be reproduced correctly on other servers in a replication setup.

To start with an `AUTO_INCREMENT` value other than 1, set that value with `CREATE TABLE` or `ALTER TABLE`, like this:

```
mysql> ALTER TABLE tbl AUTO_INCREMENT = 100;
```

InnoDB Notes

For information about AUTO_INCREMENT usage specific to InnoDB, see Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”.

MyISAM Notes

- For MyISAM tables, you can specify AUTO_INCREMENT on a secondary column in a multiple-column index. In this case, the generated value for the AUTO_INCREMENT column is calculated as MAX(*auto_increment_column*) + 1 WHERE prefix=*given-prefix*. This is useful when you want to put data into ordered groups.

  ```
  CREATE TABLE animals (
      grp ENUM('fish','mammal','bird') NOT NULL,
      id MEDIUMINT NOT NULL AUTO_INCREMENT,
      name CHAR(30) NOT NULL,
      PRIMARY KEY (grp,id)
  ) ENGINE=MyISAM;

  INSERT INTO animals (grp,name) VALUES
      ('mammal','dog'),('mammal','cat'),
      ('bird','penguin'),('fish','lax'),('mammal','whale'),
      ('bird','ostrich');

  SELECT * FROM animals ORDER BY grp,id;
  ```

  Which returns:

  ```
  +--------+----+---------+
  | grp    | id | name    |
  +--------+----+---------+
  | fish   |  1 | lax     |
  | mammal |  1 | dog     |
  | mammal |  2 | cat     |
  | mammal |  3 | whale   |
  | bird   |  1 | penguin |
  | bird   |  2 | ostrich |
  +--------+----+---------+
  ```

  In this case (when the AUTO_INCREMENT column is part of a multiple-column index), AUTO_INCREMENT values are reused if you delete the row with the biggest AUTO_INCREMENT value in any group. This happens even for MyISAM tables, for which AUTO_INCREMENT values normally are not reused.

- If the AUTO_INCREMENT column is part of multiple indexes, MySQL generates sequence values using the index that begins with the AUTO_INCREMENT column, if there is one. For example, if the animals table contained indexes PRIMARY KEY (grp, id) and INDEX (id), MySQL would ignore the PRIMARY KEY for generating sequence values. As a result, the table would contain a single sequence, not a sequence per grp value.

Further Reading

More information about AUTO_INCREMENT is available here:

- How to assign the AUTO_INCREMENT attribute to a column: Section 13.1.15, “CREATE TABLE Syntax”, and Section 13.1.7, “ALTER TABLE Syntax”.
- How AUTO_INCREMENT behaves depending on the NO_AUTO_VALUE_ON_ZERO SQL mode: Section 5.1.8, “Server SQL Modes”.
- How to use the LAST_INSERT_ID() function to find the row that contains the most recent AUTO_INCREMENT value: Section 12.14, “Information Functions”.

- Setting the `AUTO_INCREMENT` value to be used: Section 5.1.5, “Server System Variables”.
- Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”
- `AUTO_INCREMENT` and replication: Section 18.4.1.1, “Replication and AUTO_INCREMENT”.
- Server-system variables related to `AUTO_INCREMENT` (`auto_increment_increment` and `auto_increment_offset`) that can be used for replication: Section 5.1.5, “Server System Variables”.

3.7 Using MySQL with Apache

There are programs that let you authenticate your users from a MySQL database and also let you write your log files into a MySQL table.

You can change the Apache logging format to be easily readable by MySQL by putting the following into the Apache configuration file:

```
LogFormat \
        "\"%h\",%{%Y%m%d%H%M%S}t,%>s,\"%b\",\"%{Content-Type}o\",  \
        \"%U\",\"%{Referer}i\",\"%{User-Agent}i\""
```

To load a log file in that format into MySQL, you can use a statement something like this:

```
LOAD DATA INFILE '/local/access_log' INTO TABLE tbl_name
FIELDS TERMINATED BY ',' OPTIONALLY ENCLOSED BY '"' ESCAPED BY '\\'
```

The named table should be created to have columns that correspond to those that the `LogFormat` line writes to the log file.

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This chapter provides a brief overview of the MySQL command-line programs provided by Oracle Corporation. It also discusses the general syntax for specifying options when you run these programs. Most programs have options that are specific to their own operation, but the option syntax is similar for all of them. Finally, the chapter provides more detailed descriptions of individual programs, including which options they recognize.

4.1 Overview of MySQL Programs

There are many different programs in a MySQL installation. This section provides a brief overview of them. Later sections provide a more detailed description of each one. Each program's description indicates its invocation syntax and the options that it supports.

Most MySQL distributions include all of these programs, except for those programs that are platform-specific. (For example, the server startup scripts are not used on Windows.) The exception is that RPM distributions are more specialized. There is one RPM for the server, another for client programs, and so forth. If you appear to be missing one or more programs, see Chapter 2, *Installing and Upgrading MySQL*, for information on types of distributions and what they contain. It may be that you have a distribution that does not include all programs and you need to install an additional package.

Each MySQL program takes many different options. Most programs provide a `--help` option that you can use to get a description of the program's different options. For example, try `mysql --help`.

You can override default option values for MySQL programs by specifying options on the command line or in an option file. See Section 4.2, “Using MySQL Programs”, for general information on invoking programs and specifying program options.

The MySQL server, `mysqld`, is the main program that does most of the work in a MySQL installation. The server is accompanied by several related scripts that assist you in starting and stopping the server:

- `mysqld`

  The SQL daemon (that is, the MySQL server). To use client programs, `mysqld` must be running, because clients gain access to databases by connecting to the server. See Section 4.3.1, “`mysqld` — The MySQL Server”.

- `mysqld_safe`

  A server startup script. `mysqld_safe` attempts to start `mysqld`. See Section 4.3.2, “`mysqld_safe` — MySQL Server Startup Script”.

- `mysql.server`

  A server startup script. This script is used on systems that use System V-style run directories containing scripts that start system services for particular run levels. It invokes `mysqld_safe` to start the MySQL server. See Section 4.3.3, “`mysql.server` — MySQL Server Startup Script”.

- `mysqld_multi`

  A server startup script that can start or stop multiple servers installed on the system. See Section 4.3.4, “`mysqld_multi` — Manage Multiple MySQL Servers”.

Several programs perform setup operations during MySQL installation or upgrading:

- `comp_err`

  This program is used during the MySQL build/installation process. It compiles error message files from the error source files. See Section 4.4.1, “`comp_err` — Compile MySQL Error Message File”.

- `mysql_secure_installation`

  This program enables you to improve the security of your MySQL installation. See Section 4.4.2, “`mysql_secure_installation` — Improve MySQL Installation Security”.

- `mysql_ssl_rsa_setup`

  This program creates the SSL certificate and key files and RSA key-pair files required to support secure connections, if those files are missing. Files created by `mysql_ssl_rsa_setup` can be used for secure connections using SSL or RSA. See Section 4.4.3, “`mysql_ssl_rsa_setup` — Create SSL/RSA Files”.

- `mysql_tzinfo_to_sql`

  This program loads the time zone tables in the `mysql` database using the contents of the host system *zoneinfo* database (the set of files describing time zones). See Section 4.4.4, “`mysql_tzinfo_to_sql` — Load the Time Zone Tables”.

- `mysql_upgrade`

  This program is used after a MySQL upgrade operation. It checks tables for incompatibilities and repairs them if necessary, and updates the grant tables with any changes that have been made in newer versions of MySQL. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

MySQL client programs that connect to the MySQL server:

- `mysql`

  The command-line tool for interactively entering SQL statements or executing them from a file in batch mode. See Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”.

- `mysqladmin`

  A client that performs administrative operations, such as creating or dropping databases, reloading the grant tables, flushing tables to disk, and reopening log files. `mysqladmin` can also be used to retrieve version, process, and status information from the server. See Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”.

- `mysqlcheck`

  A table-maintenance client that checks, repairs, analyzes, and optimizes tables. See Section 4.5.3, “`mysqlcheck` — A Table Maintenance Program”.

- `mysqldump`

  A client that dumps a MySQL database into a file as SQL, text, or XML. See Section 4.5.4, “`mysqldump` — A Database Backup Program”.

- `mysqlimport`

  A client that imports text files into their respective tables using `LOAD DATA INFILE`. See Section 4.5.5, “`mysqlimport` — A Data Import Program”.

- `mysqlpump`

  A client that dumps a MySQL database into a file as SQL. See Section 4.5.6, “`mysqlpump` — A Database Backup Program”.

- `mysqlshow`

  A client that displays information about databases, tables, columns, and indexes. See Section 4.5.7, “`mysqlshow` — Display Database, Table, and Column Information”.

- `mysqlslap`

A client that is designed to emulate client load for a MySQL server and report the timing of each stage. It works as if multiple clients are accessing the server. See Section 4.5.8, “mysqlslap — Load Emulation Client”.

MySQL administrative and utility programs:

- innochecksum

  An offline InnoDB offline file checksum utility. See Section 4.6.2, “innochecksum — Offline InnoDB File Checksum Utility”.

- myisam_ftdump

  A utility that displays information about full-text indexes in MyISAM tables. See Section 4.6.3, “myisam_ftdump — Display Full-Text Index information”.

- myisamchk

  A utility to describe, check, optimize, and repair MyISAM tables. See Section 4.6.4, “myisamchk — MyISAM Table-Maintenance Utility”.

- myisamlog

  A utility that processes the contents of a MyISAM log file. See Section 4.6.5, “myisamlog — Display MyISAM Log File Contents”.

- myisampack

  A utility that compresses MyISAM tables to produce smaller read-only tables. See Section 4.6.6, “myisampack — Generate Compressed, Read-Only MyISAM Tables”.

- mysql_config_editor

  A utility that enables you to store authentication credentials in a secure, encrypted login path file named .mylogin.cnf. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

- mysqlbinlog

  A utility for reading statements from a binary log. The log of executed statements contained in the binary log files can be used to help recover from a crash. See Section 4.6.8, “mysqlbinlog — Utility for Processing Binary Log Files”.

- mysqldumpslow

  A utility to read and summarize the contents of a slow query log. See Section 4.6.9, “mysqldumpslow — Summarize Slow Query Log Files”.

MySQL program-development utilities:

- mysql_config

  A shell script that produces the option values needed when compiling MySQL programs. See Section 4.7.1, “mysql_config — Display Options for Compiling Clients”.

- my_print_defaults

  A utility that shows which options are present in option groups of option files. See Section 4.7.2, “my_print_defaults — Display Options from Option Files”.

- resolve_stack_dump

A utility program that resolves a numeric stack trace dump to symbols. See Section 4.7.3, “`resolve_stack_dump` — Resolve Numeric Stack Trace Dump to Symbols”.

Miscellaneous utilities:

- `lz4_decompress`

  A utility that decompresses `mysqlpump` output that was created using LZ4 compression. See Section 4.8.1, “`lz4_decompress` — Decompress mysqlpump LZ4-Compressed Output”.

- `perror`

  A utility that displays the meaning of system or MySQL error codes. See Section 4.8.2, “`perror` — Explain Error Codes”.

- `resolveip`

  A utility program that resolves a host name to an IP address or vice versa. See Section 4.8.3, “`resolveip` — Resolve Host name to IP Address or Vice Versa”.

- `zlib_decompress`

  A utility that decompresses `mysqlpump` output that was created using ZLIB compression. See Section 4.8.4, “`zlib_decompress` — Decompress mysqlpump ZLIB-Compressed Output”.

Oracle Corporation also provides the MySQL Workbench GUI tool, which is used to administer MySQL servers and databases, to create, execute, and evaluate queries, and to migrate schemas and data from other relational database management systems for use with MySQL. Additional GUI tools include MySQL Notifier and MySQL for Excel.

MySQL client programs that communicate with the server using the MySQL client/server library use the following environment variables.

Environment Variable	Meaning
`MYSQL_UNIX_PORT`	The default Unix socket file; used for connections to `localhost`
`MYSQL_TCP_PORT`	The default port number; used for TCP/IP connections
`MYSQL_PWD`	The default password
`MYSQL_DEBUG`	Debug trace options when debugging
`TMPDIR`	The directory where temporary tables and files are created

For a full list of environment variables used by MySQL programs, see Section 4.9, “MySQL Program Environment Variables”.

Use of `MYSQL_PWD` is insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”.

4.2 Using MySQL Programs

4.2.1 Invoking MySQL Programs

To invoke a MySQL program from the command line (that is, from your shell or command prompt), enter the program name followed by any options or other arguments needed to instruct the program what you want it to do. The following commands show some sample program invocations. `shell>` represents the prompt for your command interpreter; it is not part of what you type. The particular prompt you see depends on your command interpreter. Typical prompts are `$` for `sh`, `ksh`, or `bash`, `%` for `csh` or `tcsh`, and `C:\>` for the Windows `command.com` or `cmd.exe` command interpreters.

```
shell> mysql --user=root test
shell> mysqladmin extended-status variables
shell> mysqlshow --help
shell> mysqldump -u root personnel
```

Arguments that begin with a single or double dash (`-`, `--`) specify program options. Options typically indicate the type of connection a program should make to the server or affect its operational mode. Option syntax is described in Section 4.2.3, “Specifying Program Options”.

Nonoption arguments (arguments with no leading dash) provide additional information to the program. For example, the `mysql` program interprets the first nonoption argument as a database name, so the command `mysql --user=root test` indicates that you want to use the `test` database.

Later sections that describe individual programs indicate which options a program supports and describe the meaning of any additional nonoption arguments.

Some options are common to a number of programs. The most frequently used of these are the `--host` (or `-h`), `--user` (or `-u`), and `--password` (or `-p`) options that specify connection parameters. They indicate the host where the MySQL server is running, and the user name and password of your MySQL account. All MySQL client programs understand these options; they enable you to specify which server to connect to and the account to use on that server. Other connection options are `--port` (or `-P`) to specify a TCP/IP port number and `--socket` (or `-S`) to specify a Unix socket file on Unix (or named pipe name on Windows). For more information on options that specify connection options, see Section 4.2.2, “Connecting to the MySQL Server”.

You may find it necessary to invoke MySQL programs using the path name to the `bin` directory in which they are installed. This is likely to be the case if you get a “program not found” error whenever you attempt to run a MySQL program from any directory other than the `bin` directory. To make it more convenient to use MySQL, you can add the path name of the `bin` directory to your `PATH` environment variable setting. That enables you to run a program by typing only its name, not its entire path name. For example, if `mysql` is installed in `/usr/local/mysql/bin`, you can run the program by invoking it as `mysql`, and it is not necessary to invoke it as `/usr/local/mysql/bin/mysql`.

Consult the documentation for your command interpreter for instructions on setting your `PATH` variable. The syntax for setting environment variables is interpreter-specific. (Some information is given in Section 4.2.10, “Setting Environment Variables”.) After modifying your `PATH` setting, open a new console window on Windows or log in again on Unix so that the setting goes into effect.

4.2.2 Connecting to the MySQL Server

This section describes how to establish a connection to the MySQL server. For additional information if you are unable to connect, see Section 6.2.9, “Troubleshooting Problems Connecting to MySQL”.

For a client program to be able to connect to the MySQL server, it must use the proper connection parameters, such as the name of the host where the server is running and the user name and password of your MySQL account. Each connection parameter has a default value, but you can override them as necessary using program options specified either on the command line or in an option file.

The examples here use the `mysql` client program, but the principles apply to other clients such as `mysqldump`, `mysqladmin`, or `mysqlshow`.

This command invokes `mysql` without specifying any connection parameters explicitly:

```
shell> mysql
```

Because there are no parameter options, the default values apply:

- The default host name is `localhost`. On Unix, this has a special meaning, as described later.

- The default user name is `ODBC` on Windows or your Unix login name on Unix.
- No password is sent if neither `-p` nor `--password` is given.
- For `mysql`, the first nonoption argument is taken as the name of the default database. If there is no such option, `mysql` does not select a default database.

To specify the host name and user name explicitly, as well as a password, supply appropriate options on the command line:

```
shell> mysql --host=localhost --user=myname --password=mypass mydb
shell> mysql -h localhost -u myname -pmypass mydb
```

For password options, the password value is optional:

- If you use a `-p` or `--password` option and specify the password value, there must be *no space* between `-p` or `--password=` and the password following it.
- If you use a `-p` or `--password` option but do not specify the password value, the client program prompts you to enter the password. The password is not displayed as you enter it. This is more secure than giving the password on the command line. Other users on your system may be able to see a password specified on the command line by executing a command such as `ps auxw`. See Section 6.1.2.1, “End-User Guidelines for Password Security”.

As just mentioned, including the password value on the command line can be a security risk. To avoid this problem, specify the `--password` or `-p` option without any following password value:

```
shell> mysql --host=localhost --user=myname --password mydb
shell> mysql -h localhost -u myname -p mydb
```

When the password option has no password value, the client program prints a prompt and waits for you to enter the password. (In these examples, `mydb` is *not* interpreted as a password because it is separated from the preceding password option by a space.)

On some systems, the library routine that MySQL uses to prompt for a password automatically limits the password to eight characters. That is a problem with the system library, not with MySQL. Internally, MySQL does not have any limit for the length of the password. To work around the problem, change your MySQL password to a value that is eight or fewer characters long, or put your password in an option file.

On Unix, MySQL programs treat the host name `localhost` specially, in a way that is likely different from what you expect compared to other network-based programs. For connections to `localhost`, MySQL programs attempt to connect to the local server by using a Unix socket file. This occurs even if a `--port` or `-P` option is given to specify a port number. To ensure that the client makes a TCP/IP connection to the local server, use `--host` or `-h` to specify a host name value of `127.0.0.1`, or the IP address or name of the local server. You can also specify the connection protocol explicitly, even for `localhost`, by using the `--protocol=TCP` option. For example:

```
shell> mysql --host=127.0.0.1
shell> mysql --protocol=TCP
```

The `--protocol` option enables you to establish a particular type of connection even when the other options would normally default to some other protocol.

If the server is configured to accept IPv6 connections, clients can connect over IPv6 using `--host=::1`. See Section 5.1.9, “IPv6 Support”.

On Windows, you can force a MySQL client to use a named-pipe connection by specifying the `--pipe` or `--protocol=PIPE` option, or by specifying `.` (period) as the host name. If named-pipe connections

are not enabled, an error occurs. Use the `--socket` option to specify the name of the pipe if you do not want to use the default pipe name.

Connections to remote servers always use TCP/IP. This command connects to the server running on `remote.example.com` using the default port number (3306):

```
shell> mysql --host=remote.example.com
```

To specify a port number explicitly, use the `--port` or `-P` option:

```
shell> mysql --host=remote.example.com --port=13306
```

You can specify a port number for connections to a local server, too. However, as indicated previously, connections to `localhost` on Unix will use a socket file by default. You will need to force a TCP/IP connection as already described or any option that specifies a port number will be ignored.

For this command, the program uses a socket file on Unix and the `--port` option is ignored:

```
shell> mysql --port=13306 --host=localhost
```

To cause the port number to be used, invoke the program in either of these ways:

```
shell> mysql --port=13306 --host=127.0.0.1
shell> mysql --port=13306 --protocol=TCP
```

The following list summarizes the options that can be used to control how client programs connect to the server:

- `--host=`*`host_name`*, `-h` *`host_name`*

  The host where the server is running. The default value is `localhost`.

- `--password[=`*`pass_val`*`]`, `-p[`*`pass_val`*`]`

  The password of the MySQL account. As described earlier, the password value is optional, but if given, there must be *no space* between `-p` or `--password=` and the password following it. The default is to send no password.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. The server must be started with the `--enable-named-pipe` option to enable named-pipe connections.

- `--port=`*`port_num`*, `-P` *`port_num`*

  The port number to use for the connection, for connections made using TCP/IP. The default port number is 3306.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  This option explicitly specifies a protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For example, connections on Unix to `localhost` are made using a Unix socket file by default:

  ```
  shell> mysql --host=localhost
  ```

  To force a TCP/IP connection to be used instead, specify a `--protocol` option:

```
shell> mysql --host=localhost --protocol=TCP
```

The following table shows the permissible `--protocol` option values and indicates the platforms on which each value may be used. The values are not case sensitive.

`--protocol` Value	Connection Protocol	Permissible Operating Systems
`TCP`	TCP/IP connection to local or remote server	All
`SOCKET`	Unix socket file connection to local server	Unix only
`PIPE`	Named-pipe connection to local or remote server	Windows only
`MEMORY`	Shared-memory connection to local server	Windows only

- `--shared-memory-base-name=`*`name`*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--socket=`*`file_name`*, `-S` *`file_name`*

  On Unix, the name of the Unix socket file to use, for connections made using a named pipe to a local server. The default Unix socket file name is `/tmp/mysql.sock`.

  On Windows, the name of the named pipe to use, for connections to a local server. The default Windows pipe name is `MySQL`. The pipe name is not case sensitive.

  The server must be started with the `--enable-named-pipe` option to enable named-pipe connections.

- `--ssl*`

  Options that begin with `--ssl` are used for establishing a secure connection to the server using SSL, if the server is configured with SSL support. For details, see Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=`*`protocol_list`*

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=`*`user_name`*, `-u` *`user_name`*

  The user name of the MySQL account you want to use. The default user name is `ODBC` on Windows or your Unix login name on Unix.

It is possible to specify different default values to be used when you make a connection so that you need not enter them on the command line each time you invoke a client program. This can be done in a couple of ways:

- You can specify connection parameters in the `[client]` section of an option file. The relevant section of the file might look like this:

```
[client]
host=host_name
user=user_name
password=your_pass
```

Section 4.2.6, “Using Option Files”, discusses option files further.

- You can specify some connection parameters using environment variables. The host can be specified for `mysql` using `MYSQL_HOST`. The MySQL user name can be specified using `USER` (this is for Windows only). The password can be specified using `MYSQL_PWD`, although this is insecure; see Section 6.1.2.1, “End-User Guidelines for Password Security”. For a list of variables, see Section 4.9, “MySQL Program Environment Variables”.

4.2.3 Specifying Program Options

There are several ways to specify options for MySQL programs:

- List the options on the command line following the program name. This is common for options that apply to a specific invocation of the program.
- List the options in an option file that the program reads when it starts. This is common for options that you want the program to use each time it runs.
- List the options in environment variables (see Section 4.2.10, “Setting Environment Variables”). This method is useful for options that you want to apply each time the program runs. In practice, option files are used more commonly for this purpose, but Section 5.7.3, “Running Multiple MySQL Instances on Unix”, discusses one situation in which environment variables can be very helpful. It describes a handy technique that uses such variables to specify the TCP/IP port number and Unix socket file for the server and for client programs.

Options are processed in order, so if an option is specified multiple times, the last occurrence takes precedence. The following command causes `mysql` to connect to the server running on `localhost`:

```
shell> mysql -h example.com -h localhost
```

If conflicting or related options are given, later options take precedence over earlier options. The following command runs `mysql` in “no column names” mode:

```
shell> mysql --column-names --skip-column-names
```

MySQL programs determine which options are given first by examining environment variables, then by processing option files, and then by checking the command line. This means that environment variables have the lowest precedence and command-line options the highest.

For the server, one exception applies: The `mysqld-auto.cnf` option file in the data directory is processed last, so it takes precedence even over command-line options.

You can take advantage of the way that MySQL programs process options by specifying default option values for a program in an option file. That enables you to avoid typing them each time you run the program while enabling you to override the defaults if necessary by using command-line options.

4.2.4 Using Options on the Command Line

Program options specified on the command line follow these rules:

- Options are given after the command name.
- An option argument begins with one dash or two dashes, depending on whether it is a short form or long form of the option name. Many options have both short and long forms. For example, `-?` and `--help` are the short and long forms of the option that instructs a MySQL program to display its help message.
- Option names are case sensitive. `-v` and `-V` are both legal and have different meanings. (They are the corresponding short forms of the `--verbose` and `--version` options.)

- Some options take a value following the option name. For example, `-h localhost` or `--host=localhost` indicate the MySQL server host to a client program. The option value tells the program the name of the host where the MySQL server is running.
- For a long option that takes a value, separate the option name and the value by an `=` sign. For a short option that takes a value, the option value can immediately follow the option letter, or there can be a space between: `-hlocalhost` and `-h localhost` are equivalent. An exception to this rule is the option for specifying your MySQL password. This option can be given in long form as `--password=pass_val` or as `--password`. In the latter case (with no password value given), the program prompts you for the password. The password option also may be given in short form as `-ppass_val` or as `-p`. However, for the short form, if the password value is given, it must follow the option letter with *no intervening space*. The reason for this is that if a space follows the option letter, the program has no way to tell whether a following argument is supposed to be the password value or some other kind of argument. Consequently, the following two commands have two completely different meanings:

  ```
  shell> mysql -ptest
  shell> mysql -p test
  ```

  The first command instructs `mysql` to use a password value of `test`, but specifies no default database. The second instructs `mysql` to prompt for the password value and to use `test` as the default database.
- Within option names, dash (`-`) and underscore (_) may be used interchangeably. For example, `--skip-grant-tables` and `--skip_grant_tables` are equivalent. (However, the leading dashes cannot be given as underscores.)
- For options that take a numeric value, the value can be given with a suffix of `K`, `M`, or `G` (either uppercase or lowercase) to indicate a multiplier of 1024, 1024^2 or 1024^3. For example, the following command tells `mysqladmin` to ping the server 1024 times, sleeping 10 seconds between each ping:

  ```
  shell> mysqladmin --count=1K --sleep=10 ping
  ```
- When specifying file names as option values, avoid the use of the `~` shell metacharacter because it might not be interpreted as you expect.

Option values that contain spaces must be quoted when given on the command line. For example, the `--execute` (or `-e`) option can be used with `mysql` to pass SQL statements to the server. When this option is used, `mysql` executes the statements in the option value and exits. The statements must be enclosed by quotation marks. For example, you can use the following command to obtain a list of user accounts:

```
shell> mysql -u root -p --execute="SELECT User, Host FROM mysql.user"
Enter password: ******
+------+-----------+
| User | Host      |
+------+-----------+
|      | gigan     |
| root | gigan     |
|      | localhost |
| jon  | localhost |
| root | localhost |
+------+-----------+
shell>
```

Note

The long form (`--execute`) is followed by an equals sign (`=`).

If you wish to use quoted values within a statement, you will either need to escape the inner quotation marks, or use a different type of quotation marks within the statement from those used to quote the

statement itself. The capabilities of your command processor dictate your choices for whether you can use single or double quotation marks and the syntax for escaping quote characters. For example, if your command processor supports quoting with single or double quotation marks, you can use double quotation marks around the statement, and single quotation marks for any quoted values within the statement.

Multiple SQL statements may be passed in the option value on the command line, separated by semicolons:

```
shell> mysql -u root -p -e "SELECT VERSION();SELECT NOW()"
Enter password: ******
+---------------------+
| VERSION()           |
+---------------------+
| 5.8.0-m17-debug-log |
+---------------------+
+---------------------+
| NOW()               |
+---------------------+
| 2015-11-05 20:00:20 |
+---------------------+
```

4.2.5 Program Option Modifiers

Some options are “boolean” and control behavior that can be turned on or off. For example, the `mysql` client supports a `--column-names` option that determines whether or not to display a row of column names at the beginning of query results. By default, this option is enabled. However, you may want to disable it in some instances, such as when sending the output of `mysql` into another program that expects to see only data and not an initial header line.

To disable column names, you can specify the option using any of these forms:

```
--disable-column-names
--skip-column-names
--column-names=0
```

The `--disable` and `--skip` prefixes and the `=0` suffix all have the same effect: They turn the option off.

The “enabled” form of the option may be specified in any of these ways:

```
--column-names
--enable-column-names
--column-names=1
```

The values `ON`, `TRUE`, `OFF`, and `FALSE` are also recognized for boolean options (not case sensitive).

If an option is prefixed by `--loose`, a program does not exit with an error if it does not recognize the option, but instead issues only a warning:

```
shell> mysql --loose-no-such-option
mysql: WARNING: unknown option '--loose-no-such-option'
```

The `--loose` prefix can be useful when you run programs from multiple installations of MySQL on the same machine and list options in an option file. An option that may not be recognized by all versions of a program can be given using the `--loose` prefix (or `loose` in an option file). Versions of the program that recognize the option process it normally, and versions that do not recognize it issue a warning and ignore it.

The `--maximum` prefix is available for `mysqld` only and permits a limit to be placed on how large client programs can set session system variables. To do this, use a `--maximum` prefix with the variable

name. For example, `--maximum-max_heap_table_size=32M` prevents any client from making the heap table size limit larger than 32M.

The `--maximum` prefix is intended for use with system variables that have a session value. If applied to a system variable that has only a global value, an error occurs. For example, with `--maximum-query_cache_size=4M`, the server produces this error:

```
Maximum value of 'query_cache_size' cannot be set
```

4.2.6 Using Option Files

Most MySQL programs can read startup options from option files (sometimes called configuration files). Option files provide a convenient way to specify commonly used options so that they need not be entered on the command line each time you run a program.

To determine whether a program reads option files, invoke it with the `--help` option. (For `mysqld`, use `--verbose` and `--help`.) If the program reads option files, the help message indicates which files it looks for and which option groups it recognizes.

Note

A MySQL program started with the `--no-defaults` option reads no option files other than `.mylogin.cnf`.

A server started with the `persisted_globals_load` system variable disabled does not read `mysqld-auto.cnf`.

Many option files are plain text files, created using any text editor. The exceptions are:

- The `.mylogin.cnf` file that contains login path options. This is an encrypted file created by the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”. A “login path” is an option group that permits only certain options: `host`, `user`, `password`, `port` and `socket`. Client programs specify which login path to read from `.mylogin.cnf` using the `--login-path` option.

  To specify an alternative login path file name, set the `MYSQL_TEST_LOGIN_FILE` environment variable. This variable is used by the `mysql-test-run.pl` testing utility, but also is recognized by `mysql_config_editor` and by MySQL clients such as `mysql`, `mysqladmin`, and so forth.

- The `mysqld-auto.cnf` file in the data directory. This JSON-format file contains persisted system variable settings. It is created by the server upon execution of `SET PERSIST` statements. See Section 13.7.4.1, “SET Syntax for Variable Assignment”. Management of `mysqld-auto.cnf` should be left to the server and not performed manually.

MySQL looks for option files in the order described in the following discussion and reads any that exist. If an option file you want to use does not exist, create it using the appropriate method, as just discussed.

On Windows, MySQL programs read startup options from the files shown in the following table, in the specified order (top files are read first, files read later take precedence).

Table 4.1 Option Files Read on Windows Systems

File Name	Purpose
`%PROGRAMDATA%\MySQL\MySQL Server 8.0\my.ini`, `%PROGRAMDATA%\MySQL\MySQL Server 8.0\my.cnf`	Global options

File Name	Purpose
%WINDIR%\my.ini, %WINDIR%\my.cnf	Global options
C:\my.ini, C:\my.cnf	Global options
BASEDIR\my.ini, *BASEDIR*\my.cnf	Global options
defaults-extra-file	The file specified with --defaults-extra-file, if any
%APPDATA%\MySQL\.mylogin.cnf	Login path options (clients only)
DATADIR\mysqld-auto.cnf	System variables persisted with SET PERSIST (server only)

In the preceding table, %PROGRAMDATA% represents the file system directory that contains application data for all users on the host. This path defaults to C:\ProgramData on Microsoft Windows Vista and greater, and C:\Documents and Settings\All Users\Application Data on older versions of Microsoft Windows.

%WINDIR% represents the location of your Windows directory. This is commonly C:\WINDOWS. Use the following command to determine its exact location from the value of the WINDIR environment variable:

```
C:\> echo %WINDIR%
```

%APPDATA% represents the value of the Windows application data directory. Use the following command to determine its exact location from the value of the APPDATA environment variable:

```
C:\> echo %APPDATA%
```

BASEDIR represents the MySQL base installation directory. When MySQL 8.0 has been installed using MySQL Installer, this is typically C:*PROGRAMDIR*\MySQL\MySQL 8.0 Server where *PROGRAMDIR* represents the programs directory (usually Program Files on English-language versions of Windows), See Section 2.3.3, “MySQL Installer for Windows”.

DATADIR represents the MySQL data directory. As used to find mysqld-auto.cnf, its default value is the data directory location built in when MySQL was compiled, but can be changed by --datadir specified as an option-file or command-line option processed before mysqld-auto.cnf is processed.

On Unix and Unix-like systems, MySQL programs read startup options from the files shown in the following table, in the specified order (top files are read first, files read later take precedence).

Note

On Unix platforms, MySQL ignores configuration files that are world-writable. This is intentional as a security measure.

Table 4.2 Option Files Read on Unix and Unix-Like Systems

File Name	Purpose
/etc/my.cnf	Global options
/etc/mysql/my.cnf	Global options
SYSCONFDIR/my.cnf	Global options
$MYSQL_HOME/my.cnf	Server-specific options (server only)
defaults-extra-file	The file specified with --defaults-extra-file, if any
~/.my.cnf	User-specific options

File Name	Purpose
`~/.mylogin.cnf`	User-specific login path options (clients only)
`DATADIR``/mysqld-auto.cnf`	System variables persisted with `SET PERSIST` (server only)

In the preceding table, `~` represents the current user's home directory (the value of `$HOME`).

`SYSCONFDIR` represents the directory specified with the `SYSCONFDIR` option to `CMake` when MySQL was built. By default, this is the `etc` directory located under the compiled-in installation directory.

`MYSQL_HOME` is an environment variable containing the path to the directory in which the server-specific `my.cnf` file resides. If `MYSQL_HOME` is not set and you start the server using the `mysqld_safe` program, `mysqld_safe` sets it to *`BASEDIR`*, the MySQL base installation directory.

`DATADIR` represents the MySQL data directory. As used to find `mysqld-auto.cnf`, its default value is the data directory location built in when MySQL was compiled, but can be changed by `--datadir` specified as an option-file or command-line option processed before `mysqld-auto.cnf` is processed.

If multiple instances of a given option are found, the last instance takes precedence, with one exception: For `mysqld`, the *first* instance of the `--user` option is used as a security precaution, to prevent a user specified in an option file from being overridden on the command line.

The following description of option file syntax applies to files that you edit manually. This excludes `.mylogin.cnf`, which is created using `mysql_config_editor` and is encrypted, and `mysqld-auto.cnf`, which the server creates in JSON format.

Any long option that may be given on the command line when running a MySQL program can be given in an option file as well. To get the list of available options for a program, run it with the `--help` option. (For `mysqld`, use `--verbose` and `--help`.)

The syntax for specifying options in an option file is similar to command-line syntax (see Section 4.2.4, “Using Options on the Command Line”). However, in an option file, you omit the leading two dashes from the option name and you specify only one option per line. For example, `--quick` and `--host=localhost` on the command line should be specified as `quick` and `host=localhost` on separate lines in an option file. To specify an option of the form `--loose-`*`opt_name`* in an option file, write it as `loose-`*`opt_name`*.

Empty lines in option files are ignored. Nonempty lines can take any of the following forms:

- *`#comment`*, *`;comment`*

  Comment lines start with `#` or `;`. A `#` comment can start in the middle of a line as well.

- `[`*`group`*`]`

  `group` is the name of the program or group for which you want to set options. After a group line, any option-setting lines apply to the named group until the end of the option file or another group line is given. Option group names are not case sensitive.

- *`opt_name`*

  This is equivalent to `--`*`opt_name`* on the command line.

- *`opt_name=value`*

  This is equivalent to `--`*`opt_name=value`* on the command line. In an option file, you can have spaces around the `=` character, something that is not true on the command line. The value optionally can be enclosed within single quotation marks or double quotation marks, which is useful if the value contains a `#` comment character.

Leading and trailing spaces are automatically deleted from option names and values.

You can use the escape sequences `\b`, `\t`, `\n`, `\r`, `\\`, and `\s` in option values to represent the backspace, tab, newline, carriage return, backslash, and space characters. In option files, these escaping rules apply:

- A backslash followed by a valid escape sequence character is converted to the character represented by the sequence. For example, `\s` is converted to a space.
- A backslash not followed by a valid escape sequence character remains unchanged. For example, `\S` is retained as is.

The preceding rules mean that a literal backslash can be given as `\\`, or as `\` if it is not followed by a valid escape sequence character.

The rules for escape sequences in option files differ slightly from the rules for escape sequences in string literals in SQL statements. In the latter context, if "*x*" is not a valid escape sequence character, `\x` becomes "*x*" rather than `\x`. See Section 9.1.1, "String Literals".

The escaping rules for option file values are especially pertinent for Windows path names, which use `\` as a path name separator. A separator in a Windows path name must be written as `\\` if it is followed by an escape sequence character. It can be written as `\\` or `\` if it is not. Alternatively, `/` may be used in Windows path names and will be treated as `\`. Suppose that you want to specify a base directory of `C:\Program Files\MySQL\MySQL Server 8.0` in an option file. This can be done several ways. Some examples:

```
basedir="C:\Program Files\MySQL\MySQL Server 8.0"
basedir="C:\\Program Files\\MySQL\\MySQL Server 8.0"
basedir="C:/Program Files/MySQL/MySQL Server 8.0"
basedir=C:\\Program\sFiles\\MySQL\\MySQL\sServer\s8.0
```

If an option group name is the same as a program name, options in the group apply specifically to that program. For example, the `[mysqld]` and `[mysql]` groups apply to the `mysqld` server and the `mysql` client program, respectively.

The `[client]` option group is read by all client programs provided in MySQL distributions (but *not* by `mysqld`). To understand how third-party client programs that use the C API can use option files, see the C API documentation at Section 27.7.7.50, "mysql_options()".

The `[client]` group enables you to specify options that apply to all clients. For example, `[client]` is the appropriate group to use to specify the password for connecting to the server. (But make sure that the option file is accessible only by yourself, so that other people cannot discover your password.) Be sure not to put an option in the `[client]` group unless it is recognized by *all* client programs that you use. Programs that do not understand the option quit after displaying an error message if you try to run them.

List more general option groups first and more specific groups later. For example, a `[client]` group is more general because it is read by all client programs, whereas a `[mysqldump]` group is read only by `mysqldump`. Options specified later override options specified earlier, so putting the option groups in the order `[client]`, `[mysqldump]` enables `mysqldump`-specific options to override `[client]` options.

Here is a typical global option file:

```
[client]
port=3306
socket=/tmp/mysql.sock

[mysqld]
port=3306
socket=/tmp/mysql.sock
key_buffer_size=16M
max_allowed_packet=8M

[mysqldump]
```

```
quick
```

Here is a typical user option file:

```
[client]
# The following password will be sent to all standard MySQL clients
password="my password"

[mysql]
no-auto-rehash
connect_timeout=2
```

To create option groups to be read only by `mysqld` servers from specific MySQL release series, use groups with names of `[mysqld-5.7]`, `[mysqld-8.0]`, and so forth. The following group indicates that the `sql_mode` setting should be used only by MySQL servers with 8.0.x version numbers:

```
[mysqld-8.0]
sql_mode=TRADITIONAL
```

It is possible to use `!include` directives in option files to include other option files and `!includedir` to search specific directories for option files. For example, to include the `/home/mydir/myopt.cnf` file, use the following directive:

```
!include /home/mydir/myopt.cnf
```

To search the `/home/mydir` directory and read option files found there, use this directive:

```
!includedir /home/mydir
```

MySQL makes no guarantee about the order in which option files in the directory will be read.

Note

Any files to be found and included using the `!includedir` directive on Unix operating systems *must* have file names ending in `.cnf`. On Windows, this directive checks for files with the `.ini` or `.cnf` extension.

Write the contents of an included option file like any other option file. That is, it should contain groups of options, each preceded by a `[group]` line that indicates the program to which the options apply.

While an included file is being processed, only those options in groups that the current program is looking for are used. Other groups are ignored. Suppose that a `my.cnf` file contains this line:

```
!include /home/mydir/myopt.cnf
```

And suppose that `/home/mydir/myopt.cnf` looks like this:

```
[mysqladmin]
force

[mysqld]
key_buffer_size=16M
```

If `my.cnf` is processed by `mysqld`, only the `[mysqld]` group in `/home/mydir/myopt.cnf` is used. If the file is processed by `mysqladmin`, only the `[mysqladmin]` group is used. If the file is processed by any other program, no options in `/home/mydir/myopt.cnf` are used.

The `!includedir` directive is processed similarly except that all option files in the named directory are read.

If an option file contains `!include` or `!includedir` directives, files named by those directives are processed whenever the option file is processed, no matter where they appear in the file.

4.2.7 Command-Line Options that Affect Option-File Handling

Most MySQL programs that support option files handle the following options. Because these options affect option-file handling, they must be given on the command line and not in an option file. To work properly, each of these options must be given before other options, with these exceptions:

- `--print-defaults` may be used immediately after `--defaults-file`, `--defaults-extra-file`, or `--login-path`.
- On Windows, if the server is started with the `--defaults-file` and `--install` options, `--install` must be first. See Section 2.3.5.8, “Starting MySQL as a Windows Service”.

When specifying file names as option values, avoid the use of the ~ shell metacharacter because it might not be interpreted as you expect.

- `--defaults-extra-file=`*`file_name`*

  Read this option file after the global option file but (on Unix) before the user option file and (on all platforms) before the login path file. (For information about the order in which option files are used, see Section 4.2.6, “Using Option Files”.) If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=`*`file_name`*

  Read only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exceptions: Even with `--defaults-file`, `mysqld` reads `mysqld-auto.cnf` and client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=`*`str`*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, the `mysql` client normally reads the `[client]` and `[mysql]` groups. If the `--defaults-group-suffix=_other` option is given, `mysql` also reads the `[client_other]` and `[mysql_other]` groups.

- `--login-path=`*`name`*

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

  A client program reads the option group corresponding to the named login path, in addition to option groups that the program reads by default. Consider this command:

  ```
  shell> mysql --login-path=mypath
  ```

  By default, the `mysql` client reads the `[client]` and `[mysql]` option groups. So for the command shown, `mysql` reads `[client]` and `[mysql]` from other option files, and `[client]`, `[mysql]`, and `[mypath]` from the login path file.

  Client programs read the login path file even when the `--no-defaults` option is used.

  To specify an alternate login path file name, set the `MYSQL_TEST_LOGIN_FILE` environment variable.

- `--no-defaults`

Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

The exception is that client programs read the `.mylogin.cnf` login path file, if it exists, even when `--no-defaults` is used. This permits passwords to be specified in a safer way than on the command line even if `--no-defaults` is present. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.)

- `--print-defaults`

  Print the program name and all options that it gets from option files. Password values are masked.

4.2.8 Using Options to Set Program Variables

Many MySQL programs have internal variables that can be set at runtime using the `SET` statement. See Section 13.7.4.1, “SET Syntax for Variable Assignment”, and Section 5.1.6, “Using System Variables”.

Most of these program variables also can be set at server startup by using the same syntax that applies to specifying program options. For example, `mysql` has a `max_allowed_packet` variable that controls the maximum size of its communication buffer. To set the `max_allowed_packet` variable for `mysql` to a value of 16MB, use either of the following commands:

```
shell> mysql --max_allowed_packet=16777216
shell> mysql --max_allowed_packet=16M
```

The first command specifies the value in bytes. The second specifies the value in megabytes. For variables that take a numeric value, the value can be given with a suffix of `K`, `M`, or `G` (either uppercase or lowercase) to indicate a multiplier of 1024, 1024^2 or 1024^3. (For example, when used to set `max_allowed_packet`, the suffixes indicate units of kilobytes, megabytes, or gigabytes.)

In an option file, variable settings are given without the leading dashes:

```
[mysql]
max_allowed_packet=16777216
```

Or:

```
[mysql]
max_allowed_packet=16M
```

If you like, underscores in a variable name can be specified as dashes. The following option groups are equivalent. Both set the size of the server's key buffer to 512MB:

```
[mysqld]
key_buffer_size=512M

[mysqld]
key-buffer-size=512M
```

A variable can be specified by writing it in full or as any unambiguous prefix. For example, the `max_allowed_packet` variable can be set for `mysql` as `--max_a`, but not as `--max` because the latter is ambiguous:

```
shell> mysql --max=1000000
mysql: ambiguous option '--max=1000000' (max_allowed_packet, max_join_size)
```

Be aware that the use of variable prefixes can cause problems in the event that new variables are implemented for a program. A prefix that is unambiguous now might become ambiguous in the future.

Suffixes for specifying a value multiplier can be used when setting a variable at server startup, but not to set the value with `SET` at runtime. On the other hand, with `SET`, you can assign a variable's value using an expression, which is not true when you set a variable at server startup. For example, the first of the following lines is legal at server startup, but the second is not:

```
shell> mysql --max_allowed_packet=16M
shell> mysql --max_allowed_packet=16*1024*1024
```

Conversely, the second of the following lines is legal at runtime, but the first is not:

```
mysql> SET GLOBAL max_allowed_packet=16M;
mysql> SET GLOBAL max_allowed_packet=16*1024*1024;
```

4.2.9 Option Defaults, Options Expecting Values, and the = Sign

By convention, long forms of options that assign a value are written with an equals (`=`) sign, like this:

```
shell> mysql --host=tonfisk --user=jon
```

For options that require a value (that is, not having a default value), the equals sign is not required, and so the following is also valid:

```
shell> mysql --host tonfisk --user jon
```

In both cases, the `mysql` client attempts to connect to a MySQL server running on the host named “tonfisk” using an account with the user name “jon”.

Due to this behavior, problems can occasionally arise when no value is provided for an option that expects one. Consider the following example, where a user connects to a MySQL server running on host `tonfisk` as user `jon`:

```
shell> mysql --host 85.224.35.45 --user jon
Welcome to the MySQL monitor.  Commands end with ; or \g.
Your MySQL connection id is 3
Server version: 8.0.3 Source distribution

Type 'help;' or '\h' for help. Type '\c' to clear the buffer.

mysql> SELECT CURRENT_USER();
+----------------+
| CURRENT_USER() |
+----------------+
| jon@%          |
+----------------+
1 row in set (0.00 sec)
```

Omitting the required value for one of these option yields an error, such as the one shown here:

```
shell> mysql --host 85.224.35.45 --user
mysql: option '--user' requires an argument
```

In this case, `mysql` was unable to find a value following the `--user` option because nothing came after it on the command line. However, if you omit the value for an option that is *not* the last option to be used, you obtain a different error that you may not be expecting:

```
shell> mysql --host --user jon
```

```
ERROR 2005 (HY000): Unknown MySQL server host '--user' (1)
```

Because `mysql` assumes that any string following `--host` on the command line is a host name, `--host --user` is interpreted as `--host=--user`, and the client attempts to connect to a MySQL server running on a host named “--user”.

Options having default values always require an equals sign when assigning a value; failing to do so causes an error. For example, the MySQL server `--log-error` option has the default value *`host_name`*`.err`, where *`host_name`* is the name of the host on which MySQL is running. Assume that you are running MySQL on a computer whose host name is “tonfisk”, and consider the following invocation of `mysqld_safe`:

```
shell> mysqld_safe &
[1] 11699
shell> 080112 12:53:40 mysqld_safe Logging to '/usr/local/mysql/var/tonfisk.err'.
080112 12:53:40 mysqld_safe Starting mysqld daemon with databases from /usr/local/mysql/var
shell>
```

After shutting down the server, restart it as follows:

```
shell> mysqld_safe --log-error &
[1] 11699
shell> 080112 12:53:40 mysqld_safe Logging to '/usr/local/mysql/var/tonfisk.err'.
080112 12:53:40 mysqld_safe Starting mysqld daemon with databases from /usr/local/mysql/var
shell>
```

The result is the same, since `--log-error` is not followed by anything else on the command line, and it supplies its own default value. (The `&` character tells the operating system to run MySQL in the background; it is ignored by MySQL itself.) Now suppose that you wish to log errors to a file named `my-errors.err`. You might try starting the server with `--log-error my-errors`, but this does not have the intended effect, as shown here:

```
shell> mysqld_safe --log-error my-errors &
[1] 31357
shell> 080111 22:53:31 mysqld_safe Logging to '/usr/local/mysql/var/tonfisk.err'.
080111 22:53:32 mysqld_safe Starting mysqld daemon with databases from /usr/local/mysql/var
080111 22:53:34 mysqld_safe mysqld from pid file /usr/local/mysql/var/tonfisk.pid ended

[1]+  Done                    ./mysqld_safe --log-error my-errors
```

The server attempted to start using `/usr/local/mysql/var/tonfisk.err` as the error log, but then shut down. Examining the last few lines of this file shows the reason:

```
shell> tail /usr/local/mysql/var/tonfisk.err
2013-09-24T15:36:22.278034Z 0 [ERROR] Too many arguments (first extra is 'my-errors').
2013-09-24T15:36:22.278059Z 0 [Note] Use --verbose --help to get a list of available options!
2013-09-24T15:36:22.278076Z 0 [ERROR] Aborting
2013-09-24T15:36:22.279704Z 0 [Note] InnoDB: Starting shutdown...
2013-09-24T15:36:23.777471Z 0 [Note] InnoDB: Shutdown completed; log sequence number 2319086
2013-09-24T15:36:23.780134Z 0 [Note] mysqld: Shutdown complete
```

Because the `--log-error` option supplies a default value, you must use an equals sign to assign a different value to it, as shown here:

```
shell> mysqld_safe --log-error=my-errors &
[1] 31437
shell> 080111 22:54:15 mysqld_safe Logging to '/usr/local/mysql/var/my-errors.err'.
080111 22:54:15 mysqld_safe Starting mysqld daemon with databases from /usr/local/mysql/var

shell>
```

Now the server has been started successfully, and is logging errors to the file `/usr/local/mysql/var/my-errors.err`.

Similar issues can arise when specifying option values in option files. For example, consider a `my.cnf` file that contains the following:

```
[mysql]

host
user
```

When the `mysql` client reads this file, these entries are parsed as `--host --user` or `--host=--user`, with the result shown here:

```
shell> mysql
ERROR 2005 (HY000): Unknown MySQL server host '--user' (1)
```

However, in option files, an equals sign is not assumed. Suppose the `my.cnf` file is as shown here:

```
[mysql]

user jon
```

Trying to start `mysql` in this case causes a different error:

```
shell> mysql
mysql: unknown option '--user jon'
```

A similar error would occur if you were to write `host tonfisk` in the option file rather than `host=tonfisk`. Instead, you must use the equals sign:

```
[mysql]

user=jon
```

Now the login attempt succeeds:

```
shell> mysql
Welcome to the MySQL monitor.  Commands end with ; or \g.
Your MySQL connection id is 5
Server version: 8.0.3 Source distribution

Type 'help;' or '\h' for help. Type '\c' to clear the buffer.

mysql> SELECT USER();
+---------------+
| USER()        |
+---------------+
| jon@localhost |
+---------------+
1 row in set (0.00 sec)
```

This is not the same behavior as with the command line, where the equals sign is not required:

```
shell> mysql --user jon --host tonfisk
Welcome to the MySQL monitor.  Commands end with ; or \g.
Your MySQL connection id is 6
Server version: 8.0.3 Source distribution

Type 'help;' or '\h' for help. Type '\c' to clear the buffer.

mysql> SELECT USER();
+---------------+
| USER()        |
+---------------+
| jon@tonfisk   |
```

```
+---------------+
1 row in set (0.00 sec)
```

Specifying an option requiring a value without a value in an option file causes the server to abort with an error. Suppose that `my.cnf` contains the following:

```
[mysqld]
log_error
relay_log
relay_log_index
```

This causes the server to fail on startup, as shown here:

```
shell> mysqld_safe &

130924 10:41:46 mysqld_safe Logging to '/home/jon/bin/mysql/var/tonfisk.err'.
130924 10:41:46 mysqld_safe Starting mysqld daemon with databases from /home/jon/bin/mysql/var
130924 10:41:47 mysqld_safe mysqld from pid file /home/jon/bin/mysql/var/tonfisk.pid ended
```

The `--log-error` option does not require an argument; however, the `--relay-log` option requires one, as shown in the error log (which in the absence of a specified value, defaults to *`datadir/hostname.err`*):

```
shell> tail -n 3 ../var/tonfisk.err

130924 10:41:46 mysqld_safe Starting mysqld daemon with databases from /home/jon/bin/mysql/var
2013-09-24T15:41:47.217180Z 0 [ERROR] /home/jon/bin/mysql/libexec/mysqld: option '--relay-log' requires
2013-09-24T15:41:47.217479Z 0 [ERROR] Aborting
```

This is a change from previous behavior, where the server would have interpreted the last two lines in the example `my.cnf` file as `--relay-log=relay_log_index` and created a relay log file using “relay_log_index” as the base name. (Bug #25192)

4.2.10 Setting Environment Variables

Environment variables can be set at the command prompt to affect the current invocation of your command processor, or set permanently to affect future invocations. To set a variable permanently, you can set it in a startup file or by using the interface provided by your system for this purpose. Consult the documentation for your command interpreter for specific details. Section 4.9, “MySQL Program Environment Variables”, lists all environment variables that affect MySQL program operation.

To specify a value for an environment variable, use the syntax appropriate for your command processor. For example, on Windows, you can set the `USER` variable to specify your MySQL account name. To do so, use this syntax:

```
SET USER=your_name
```

The syntax on Unix depends on your shell. Suppose that you want to specify the TCP/IP port number using the `MYSQL_TCP_PORT` variable. Typical syntax (such as for `sh`, `ksh`, `bash`, `zsh`, and so on) is as follows:

```
MYSQL_TCP_PORT=3306
export MYSQL_TCP_PORT
```

The first command sets the variable, and the `export` command exports the variable to the shell environment so that its value becomes accessible to MySQL and other processes.

For `csh` and `tcsh`, use `setenv` to make the shell variable available to the environment:

```
setenv MYSQL_TCP_PORT 3306
```

The commands to set environment variables can be executed at your command prompt to take effect immediately, but the settings persist only until you log out. To have the settings take effect each time you log in, use the interface provided by your system or place the appropriate command or commands in a startup file that your command interpreter reads each time it starts.

On Windows, you can set environment variables using the System Control Panel (under Advanced).

On Unix, typical shell startup files are `.bashrc` or `.bash_profile` for `bash`, or `.tcshrc` for `tcsh`.

Suppose that your MySQL programs are installed in `/usr/local/mysql/bin` and that you want to make it easy to invoke these programs. To do this, set the value of the `PATH` environment variable to include that directory. For example, if your shell is `bash`, add the following line to your `.bashrc` file:

```
PATH=${PATH}:/usr/local/mysql/bin
```

`bash` uses different startup files for login and nonlogin shells, so you might want to add the setting to `.bashrc` for login shells and to `.bash_profile` for nonlogin shells to make sure that `PATH` is set regardless.

If your shell is `tcsh`, add the following line to your `.tcshrc` file:

```
setenv PATH ${PATH}:/usr/local/mysql/bin
```

If the appropriate startup file does not exist in your home directory, create it with a text editor.

After modifying your `PATH` setting, open a new console window on Windows or log in again on Unix so that the setting goes into effect.

4.3 MySQL Server and Server-Startup Programs

This section describes `mysqld`, the MySQL server, and several programs that are used to start the server.

4.3.1 `mysqld` — The MySQL Server

`mysqld`, also known as MySQL Server, is the main program that does most of the work in a MySQL installation. MySQL Server manages access to the MySQL data directory that contains databases and tables. The data directory is also the default location for other information such as log files and status files.

When MySQL server starts, it listens for network connections from client programs and manages access to databases on behalf of those clients.

The `mysqld` program has many options that can be specified at startup. For a complete list of options, run this command:

```
shell> mysqld --verbose --help
```

MySQL Server also has a set of system variables that affect its operation as it runs. System variables can be set at server startup, and many of them can be changed at runtime to effect dynamic server reconfiguration. MySQL Server also has a set of status variables that provide information about its operation. You can monitor these status variables to access runtime performance characteristics.

For a full description of MySQL Server command options, system variables, and status variables, see Section 5.1, “The MySQL Server”. For information about installing MySQL and setting up the initial configuration, see Chapter 2, *Installing and Upgrading MySQL*.

4.3.2 mysqld_safe — MySQL Server Startup Script

mysqld_safe is the recommended way to start a mysqld server on Unix. mysqld_safe adds some safety features such as restarting the server when an error occurs and logging runtime information to an error log file. A description of error logging is given later in this section.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, mysqld_safe is not installed because it is unnecessary. For more information, see Section 2.5.9, “Managing MySQL Server with systemd”.

mysqld_safe tries to start an executable named mysqld. To override the default behavior and specify explicitly the name of the server you want to run, specify a --mysqld or --mysqld-version option to mysqld_safe. You can also use --ledir to indicate the directory where mysqld_safe should look for the server.

Many of the options to mysqld_safe are the same as the options to mysqld. See Section 5.1.4, “Server Command Options”.

Options unknown to mysqld_safe are passed to mysqld if they are specified on the command line, but ignored if they are specified in the [mysqld_safe] group of an option file. See Section 4.2.6, “Using Option Files”.

mysqld_safe reads all options from the [mysqld], [server], and [mysqld_safe] sections in option files. For example, if you specify a [mysqld] section like this, mysqld_safe will find and use the --log-error option:

```
[mysqld]
log-error=error.log
```

For backward compatibility, mysqld_safe also reads [safe_mysqld] sections, but to be current you should rename such sections to [mysqld_safe].

mysqld_safe accepts options on the command line and in option files, as described in the following table. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.3 mysqld_safe Options

Format	Description
--basedir	Path to MySQL installation directory
--core-file-size	Size of core file that mysqld should be able to create
--datadir	Path to data directory
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--help	Display help message and exit
--ledir	Path to directory where server is located
--log-error	Write error log to named file
--malloc-lib	Alternative malloc library to use for mysqld
--mysqld	Name of server program to start (in ledir directory)
--mysqld-safe-log-timestamps	Timestamp format for logging
--mysqld-version	Suffix for server program name

Format	Description
--nice	Use nice program to set server scheduling priority
--no-defaults	Read no option files
--open-files-limit	Number of files that mysqld should be able to open
--pid-file	Path name of server process ID file
--plugin-dir	Directory where plugins are installed
--port	Port number on which to listen for TCP/IP connections
--skip-kill-mysqld	Do not try to kill stray mysqld processes
--skip-syslog	Do not write error messages to syslog; use error log file
--socket	Socket file on which to listen for Unix socket connections
--syslog	Write error messages to syslog
--syslog-tag	Tag suffix for messages written to syslog
--timezone	Set TZ time zone environment variable to named value
--user	Run mysqld as user having name user_name or numeric user ID user_id

- `--help`

  Display a help message and exit.

- `--basedir=`*`dir_name`*

  The path to the MySQL installation directory.

- `--core-file-size=`*`size`*

  The size of the core file that `mysqld` should be able to create. The option value is passed to `ulimit -c`.

- `--datadir=`*`dir_name`*

  The path to the data directory.

- `--defaults-extra-file=`*`file_name`*

  Read this option file in addition to the usual option files. If the file does not exist or is otherwise inaccessible, the server will exit with an error. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name. This must be the first option on the command line if it is used.

  For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--defaults-file=`*`file_name`*

  Use only the given option file. If the file does not exist or is otherwise inaccessible, the server will exit with an error. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name. This must be the first option on the command line if it is used.

  For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--ledir=`*`dir_name`*

  If `mysqld_safe` cannot find the server, use this option to indicate the path name to the directory where the server is located.

This option is accepted only on the command line, not in option files. On platforms that use systemd, the value can be specified in the value of MYSQLD_OPTS. See Section 2.5.9, “Managing MySQL Server with systemd”.

- --log-error=*file_name*

  Write the error log to the given file. See Section 5.4.2, “The Error Log”.

- --mysqld-safe-log-timestamps

  This option controls the format for timestamps in log output produced by mysqld_safe. The following list describes the permitted values. For any other value, mysqld_safe logs a warning and uses UTC format.

  - UTC, utc

    ISO 8601 UTC format (same as --log_timestamps=UTC for the server). This is the default.

  - SYSTEM, system

    ISO 8601 local time format (same as --log_timestamps=SYSTEM for the server).

  - HYPHEN, hyphen

    YY-MM-DD h:mm:ss format, as in mysqld_safe for MySQL 5.6.

  - LEGACY, legacy

    YYMMDD hh:mm:ss format, as in mysqld_safe prior to MySQL 5.6.

- --malloc-lib=[*lib_name*]

  The name of the library to use for memory allocation instead of the system malloc() library. The option value must be one of the directories /usr/lib, /usr/lib64, /usr/lib/i386-linux-gnu, or /usr/lib/x86_64-linux-gnu.

  The --malloc-lib option works by modifying the LD_PRELOAD environment value to affect dynamic linking to enable the loader to find the memory-allocation library when mysqld runs:

  - If the option is not given, or is given without a value (--malloc-lib=), LD_PRELOAD is not modified and no attempt is made to use tcmalloc.
  - If the option is given as --malloc-lib=tcmalloc, mysqld_safe looks for a tcmalloc library in /usr/lib. If tmalloc is found, its path name is added to the beginning of the LD_PRELOAD value for mysqld. If tcmalloc is not found, mysqld_safe aborts with an error.
  - If the option is given as --malloc-lib=*/path/to/some/library*, that full path is added to the beginning of the LD_PRELOAD value. If the full path points to a nonexistent or unreadable file, mysqld_safe aborts with an error.
  - For cases where mysqld_safe adds a path name to LD_PRELOAD, it adds the path to the beginning of any existing value the variable already has.

  Note

  On systems that manage the server using systemd, mysqld_safe is not available. Instead, specify the allocation library by setting LD_PRELOAD in /etc/sysconfig/mysql.

  Linux users can use the libtcmalloc_minimal.so library on any platform for which a tcmalloc package is installed in /usr/lib by adding these lines to the my.cnf file:

```
[mysqld_safe]
malloc-lib=tcmalloc
```

To use a specific tcmalloc library, specify its full path name. Example:

```
[mysqld_safe]
malloc-lib=/opt/lib/libtcmalloc_minimal.so
```

- --mysqld=*prog_name*

  The name of the server program (in the ledir directory) that you want to start. This option is needed if you use the MySQL binary distribution but have the data directory outside of the binary distribution. If mysqld_safe cannot find the server, use the --ledir option to indicate the path name to the directory where the server is located.

  This option is accepted only on the command line, not in option files. On platforms that use systemd, the value can be specified in the value of MYSQLD_OPTS. See Section 2.5.9, “Managing MySQL Server with systemd”.

- --mysqld-version=*suffix*

  This option is similar to the --mysqld option, but you specify only the suffix for the server program name. The base name is assumed to be mysqld. For example, if you use --mysqld-version=debug, mysqld_safe starts the mysqld-debug program in the ledir directory. If the argument to --mysqld-version is empty, mysqld_safe uses mysqld in the ledir directory.

  This option is accepted only on the command line, not in option files. On platforms that use systemd, the value can be specified in the value of MYSQLD_OPTS. See Section 2.5.9, “Managing MySQL Server with systemd”.

- --nice=*priority*

  Use the nice program to set the server's scheduling priority to the given value.

- --no-defaults

  Do not read any option files. If program startup fails due to reading unknown options from an option file, --no-defaults can be used to prevent them from being read. This must be the first option on the command line if it is used.

  For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- --open-files-limit=*count*

  The number of files that mysqld should be able to open. The option value is passed to ulimit -n.

  Note

  You must start mysqld_safe as root for this to function properly.

- --pid-file=*file_name*

  The path name that mysqld should use for its process ID file.

- --plugin-dir=*dir_name*

  The path name of the plugin directory.

- --port=*port_num*

The port number that the server should use when listening for TCP/IP connections. The port number must be 1024 or higher unless the server is started by the root system user.

- --skip-kill-mysqld

  Do not try to kill stray mysqld processes at startup. This option works only on Linux.

- --socket=*path*

  The Unix socket file that the server should use when listening for local connections.

- --syslog, --skip-syslog

  --syslog causes error messages to be sent to syslog on systems that support the logger program. --skip-syslog suppresses the use of syslog; messages are written to an error log file.

  When syslog is used, the daemon.err facility/severity is used for all log messages.

  Using these options to control mysqld logging is deprecated. Use the server log_syslog system variable instead. To control the facility, use the server log_syslog_facility system variable. See Section 5.4.2, “The Error Log”.

- --syslog-tag=*tag*

  For logging to syslog, messages from mysqld_safe and mysqld are written with identifiers of mysqld_safe and mysqld, respectively. To specify a suffix for the identifiers, use --syslog-tag=*tag*, which modifies the identifiers to be mysqld_safe-*tag* and mysqld-*tag*.

  Using this option to control mysqld logging is deprecated. Use the server log_syslog_tag system variable instead. See Section 5.4.2, “The Error Log”.

- --timezone=*timezone*

  Set the TZ time zone environment variable to the given option value. Consult your operating system documentation for legal time zone specification formats.

- --user={*user_name*|*user_id*}

  Run the mysqld server as the user having the name *user_name* or the numeric user ID *user_id*. (“User” in this context refers to a system login account, not a MySQL user listed in the grant tables.)

If you execute mysqld_safe with the --defaults-file or --defaults-extra-file option to name an option file, the option must be the first one given on the command line or the option file will not be used. For example, this command will not use the named option file:

```
mysql> mysqld_safe --port=port_num --defaults-file=file_name
```

Instead, use the following command:

```
mysql> mysqld_safe --defaults-file=file_name --port=port_num
```

The mysqld_safe script is written so that it normally can start a server that was installed from either a source or a binary distribution of MySQL, even though these types of distributions typically install the server in slightly different locations. (See Section 2.1.4, “Installation Layouts”.) mysqld_safe expects one of the following conditions to be true:

- The server and databases can be found relative to the working directory (the directory from which mysqld_safe is invoked). For binary distributions, mysqld_safe looks under its working directory for bin and data directories. For source distributions, it looks for libexec and var directories. This condition should be met if you execute mysqld_safe from your MySQL installation directory (for example, /usr/local/mysql for a binary distribution).

- If the server and databases cannot be found relative to the working directory, mysqld_safe attempts to locate them by absolute path names. Typical locations are /usr/local/libexec and /usr/local/var. The actual locations are determined from the values configured into the distribution at the time it was built. They should be correct if MySQL is installed in the location specified at configuration time.

Because mysqld_safe tries to find the server and databases relative to its own working directory, you can install a binary distribution of MySQL anywhere, as long as you run mysqld_safe from the MySQL installation directory:

```
shell> cd mysql_installation_directory
shell> bin/mysqld_safe &
```

If mysqld_safe fails, even when invoked from the MySQL installation directory, specify the --ledir and --datadir options to indicate the directories in which the server and databases are located on your system.

mysqld_safe tries to use the sleep and date system utilities to determine how many times per second it has attempted to start. If these utilities are present and the attempted starts per second is greater than 5, mysqld_safe waits 1 full second before starting again. This is intended to prevent excessive CPU usage in the event of repeated failures. (Bug #11761530, Bug #54035)

When you use mysqld_safe to start mysqld, mysqld_safe arranges for error (and notice) messages from itself and from mysqld to go to the same destination.

There are several mysqld_safe options for controlling the destination of these messages:

- --log-error=*file_name*: Write error messages to the named error file.
- --syslog: Write error messages to syslog on systems that support the logger program.
- --skip-syslog: Do not write error messages to syslog. Messages are written to the default error log file (*host_name*.err in the data directory), or to a named file if the --log-error option is given.

If none of these options is given, the default is --skip-syslog.

When mysqld_safe writes a message, notices go to the logging destination (syslog or the error log file) and stdout. Errors go to the logging destination and stderr.

Note

Controlling mysqld logging from mysqld_safe is deprecated. Use the server's native syslog support instead. For more information, see Section 5.4.2, “The Error Log”.

4.3.3 mysql.server — MySQL Server Startup Script

MySQL distributions on Unix and Unix-like system include a script named mysql.server, which starts the MySQL server using mysqld_safe. It can be used on systems such as Linux and Solaris that use System V-style run directories to start and stop system services. It is also used by the macOS Startup Item for MySQL.

mysql.server is the script name as used within the MySQL source tree. The installed name might be different; for example, mysqld or mysql. In the following discussion, adjust the name mysql.server as appropriate for your system.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown.

> On these platforms, mysql.server and mysqld_safe are not installed because they are unnecessary. For more information, see Section 2.5.9, “Managing MySQL Server with systemd”.

To start or stop the server manually using the mysql.server script, invoke it from the command line with start or stop arguments:

```
shell> mysql.server start
shell> mysql.server stop
```

mysql.server changes location to the MySQL installation directory, then invokes mysqld_safe. To run the server as some specific user, add an appropriate user option to the [mysqld] group of the global /etc/my.cnf option file, as shown later in this section. (It is possible that you must edit mysql.server if you've installed a binary distribution of MySQL in a nonstandard location. Modify it to change location into the proper directory before it runs mysqld_safe. If you do this, your modified version of mysql.server may be overwritten if you upgrade MySQL in the future; make a copy of your edited version that you can reinstall.)

mysql.server stop stops the server by sending a signal to it. You can also stop the server manually by executing mysqladmin shutdown.

To start and stop MySQL automatically on your server, you must add start and stop commands to the appropriate places in your /etc/rc* files:

- If you use the Linux server RPM package (MySQL-server-*VERSION*.rpm), or a native Linux package installation, the mysql.server script may be installed in the /etc/init.d directory with the name mysqld or mysql. See Section 2.5.4, “Installing MySQL on Linux Using RPM Packages from Oracle”, for more information on the Linux RPM packages.

- If you install MySQL from a source distribution or using a binary distribution format that does not install mysql.server automatically, you can install the script manually. It can be found in the support-files directory under the MySQL installation directory or in a MySQL source tree. Copy the script to the /etc/init.d directory with the name mysql and make it executable:

  ```
  shell> cp mysql.server /etc/init.d/mysql
  shell> chmod +x /etc/init.d/mysql
  ```

  After installing the script, the commands needed to activate it to run at system startup depend on your operating system. On Linux, you can use chkconfig:

  ```
  shell> chkconfig --add mysql
  ```

  On some Linux systems, the following command also seems to be necessary to fully enable the mysql script:

  ```
  shell> chkconfig --level 345 mysql on
  ```

- On FreeBSD, startup scripts generally should go in /usr/local/etc/rc.d/. Install the mysql.server script as /usr/local/etc/rc.d/mysql.server.sh to enable automatic startup. The rc(8) manual page states that scripts in this directory are executed only if their base name matches the *.sh shell file name pattern. Any other files or directories present within the directory are silently ignored.

- As an alternative to the preceding setup, some operating systems also use /etc/rc.local or /etc/init.d/boot.local to start additional services on startup. To start up MySQL using this method, append a command like the one following to the appropriate startup file:

  ```
  /bin/sh -c 'cd /usr/local/mysql; ./bin/mysqld_safe --user=mysql &'
  ```

- For other systems, consult your operating system documentation to see how to install startup scripts.

`mysql.server` reads options from the `[mysql.server]` and `[mysqld]` sections of option files. For backward compatibility, it also reads `[mysql_server]` sections, but to be current you should rename such sections to `[mysql.server]`.

You can add options for `mysql.server` in a global `/etc/my.cnf` file. A typical `my.cnf` file might look like this:

```
[mysqld]
datadir=/usr/local/mysql/var
socket=/var/tmp/mysql.sock
port=3306
user=mysql

[mysql.server]
basedir=/usr/local/mysql
```

The `mysql.server` script supports the options shown in the following table. If specified, they *must* be placed in an option file, not on the command line. `mysql.server` supports only `start` and `stop` as command-line arguments.

Table 4.4 mysql.server Option-File Options

Option Name	Description	Type
`basedir`	Path to MySQL installation directory	directory name
`datadir`	Path to MySQL data directory	directory name
`pid-file`	File in which server should write its process ID	file name
`service-startup-timeout`	How long to wait for server startup	integer

- `basedir=`*`dir_name`*

  The path to the MySQL installation directory.

- `datadir=`*`dir_name`*

  The path to the MySQL data directory.

- `pid-file=`*`file_name`*

  The path name of the file in which the server should write its process ID.

  If this option is not given, `mysql.server` uses a default value of *`host_name`*`.pid`. The PID file value passed to `mysqld_safe` overrides any value specified in the `[mysqld_safe]` option file group. Because `mysql.server` reads the `[mysqld]` option file group but not the `[mysqld_safe]` group, you can ensure that `mysqld_safe` gets the same value when invoked from `mysql.server` as when invoked manually by putting the same `pid-file` setting in both the `[mysqld_safe]` and `[mysqld]` groups.

- `service-startup-timeout=`*`seconds`*

  How long in seconds to wait for confirmation of server startup. If the server does not start within this time, `mysql.server` exits with an error. The default value is 900. A value of 0 means not to wait at all for startup. Negative values mean to wait forever (no timeout).

4.3.4 `mysqld_multi` — Manage Multiple MySQL Servers

mysqld_multi is designed to manage several mysqld processes that listen for connections on different Unix socket files and TCP/IP ports. It can start or stop servers, or report their current status.

Note

For some Linux platforms, MySQL installation from RPM or Debian packages includes systemd support for managing MySQL server startup and shutdown. On these platforms, mysqld_multi is not installed because it is unnecessary. For information about using systemd to handle multiple MySQL instances, see Section 2.5.9, “Managing MySQL Server with systemd”.

mysqld_multi searches for groups named [mysqldN] in my.cnf (or in the file named by the --defaults-file option). *N* can be any positive integer. This number is referred to in the following discussion as the option group number, or *GNR*. Group numbers distinguish option groups from one another and are used as arguments to mysqld_multi to specify which servers you want to start, stop, or obtain a status report for. Options listed in these groups are the same that you would use in the [mysqld] group used for starting mysqld. (See, for example, Section 2.9.5, “Starting and Stopping MySQL Automatically”.) However, when using multiple servers, it is necessary that each one use its own value for options such as the Unix socket file and TCP/IP port number. For more information on which options must be unique per server in a multiple-server environment, see Section 5.7, “Running Multiple MySQL Instances on One Machine”.

To invoke mysqld_multi, use the following syntax:

```
shell> mysqld_multi [options] {start|stop|reload|report} [GNR[,GNR] ...]
```

start, stop, reload (stop and restart), and report indicate which operation to perform. You can perform the designated operation for a single server or multiple servers, depending on the *GNR* list that follows the option name. If there is no list, mysqld_multi performs the operation for all servers in the option file.

Each *GNR* value represents an option group number or range of group numbers. The value should be the number at the end of the group name in the option file. For example, the *GNR* for a group named [mysqld17] is 17. To specify a range of numbers, separate the first and last numbers by a dash. The *GNR* value 10-13 represents groups [mysqld10] through [mysqld13]. Multiple groups or group ranges can be specified on the command line, separated by commas. There must be no whitespace characters (spaces or tabs) in the *GNR* list; anything after a whitespace character is ignored.

This command starts a single server using option group [mysqld17]:

```
shell> mysqld_multi start 17
```

This command stops several servers, using option groups [mysqld8] and [mysqld10] through [mysqld13]:

```
shell> mysqld_multi stop 8,10-13
```

For an example of how you might set up an option file, use this command:

```
shell> mysqld_multi --example
```

mysqld_multi searches for option files as follows:

- With --no-defaults, no option files are read.
- With --defaults-file=*file_name*, only the named file is read.
- Otherwise, option files in the standard list of locations are read, including any file named by the --defaults-extra-file=*file_name* option, if one is given. (If the option is given multiple times, the last value is used.)

Option files read are searched for [mysqld_multi] and [mysqldN] option groups. The [mysqld_multi] group can be used for options to mysqld_multi itself. [mysqldN] groups can be used for options passed to specific mysqld instances.

The [mysqld] or [mysqld_safe] groups can be used for common options read by all instances of mysqld or mysqld_safe. You can specify a --defaults-file=*file_name* option to use a different configuration file for that instance, in which case the [mysqld] or [mysqld_safe] groups from that file will be used for that instance.

mysqld_multi supports the following options.

- --help

  Display a help message and exit.

- --example

  Display a sample option file.

- --log=*file_name*

  Specify the name of the log file. If the file exists, log output is appended to it.

- --mysqladmin=*prog_name*

  The mysqladmin binary to be used to stop servers.

- --mysqld=*prog_name*

  The mysqld binary to be used. Note that you can specify mysqld_safe as the value for this option also. If you use mysqld_safe to start the server, you can include the mysqld or ledir options in the corresponding [mysqldN] option group. These options indicate the name of the server that mysqld_safe should start and the path name of the directory where the server is located. (See the descriptions for these options in Section 4.3.2, “mysqld_safe — MySQL Server Startup Script”.) Example:

  ```
  [mysqld38]
  mysqld = mysqld-debug
  ledir  = /opt/local/mysql/libexec
  ```

- --no-log

  Print log information to stdout rather than to the log file. By default, output goes to the log file.

- --password=*password*

  The password of the MySQL account to use when invoking mysqladmin. Note that the password value is not optional for this option, unlike for other MySQL programs.

- --silent

  Silent mode; disable warnings.

- --tcp-ip

  Connect to each MySQL server through the TCP/IP port instead of the Unix socket file. (If a socket file is missing, the server might still be running, but accessible only through the TCP/IP port.) By default, connections are made using the Unix socket file. This option affects stop and report operations.

- --user=*user_name*

  The user name of the MySQL account to use when invoking mysqladmin.

- `--verbose`

  Be more verbose.

- `--version`

  Display version information and exit.

Some notes about `mysqld_multi`:

- **Most important**: Before using `mysqld_multi` be sure that you understand the meanings of the options that are passed to the `mysqld` servers and *why* you would want to have separate `mysqld` processes. Beware of the dangers of using multiple `mysqld` servers with the same data directory. Use separate data directories, unless you *know* what you are doing. Starting multiple servers with the same data directory does *not* give you extra performance in a threaded system. See Section 5.7, "Running Multiple MySQL Instances on One Machine".

  Important

  Make sure that the data directory for each server is fully accessible to the Unix account that the specific `mysqld` process is started as. *Do not* use the Unix *`root`* account for this, unless you *know* what you are doing. See Section 6.1.5, "How to Run MySQL as a Normal User".

- Make sure that the MySQL account used for stopping the `mysqld` servers (with the `mysqladmin` program) has the same user name and password for each server. Also, make sure that the account has the `SHUTDOWN` privilege. If the servers that you want to manage have different user names or passwords for the administrative accounts, you might want to create an account on each server that has the same user name and password. For example, you might set up a common `multi_admin` account by executing the following commands for each server:

  ```
  shell> mysql -u root -S /tmp/mysql.sock -p
  Enter password:
  mysql> CREATE USER 'multi_admin'@'localhost' IDENTIFIED BY 'multipass';
  mysql> GRANT SHUTDOWN ON *.* TO 'multi_admin'@'localhost';
  ```

  See Section 6.2, "The MySQL Access Privilege System". You have to do this for each `mysqld` server. Change the connection parameters appropriately when connecting to each one. Note that the host name part of the account name must permit you to connect as `multi_admin` from the host where you want to run `mysqld_multi`.

- The Unix socket file and the TCP/IP port number must be different for every `mysqld`. (Alternatively, if the host has multiple network addresses, you can use `--bind-address` to cause different servers to listen to different interfaces.)

- The `--pid-file` option is very important if you are using `mysqld_safe` to start `mysqld` (for example, `--mysqld=mysqld_safe`) Every `mysqld` should have its own process ID file. The advantage of using `mysqld_safe` instead of `mysqld` is that `mysqld_safe` monitors its `mysqld` process and restarts it if the process terminates due to a signal sent using `kill -9` or for other reasons, such as a segmentation fault. Please note that the `mysqld_safe` script might require that you start it from a certain place. This means that you might have to change location to a certain directory before running `mysqld_multi`. If you have problems starting, please see the `mysqld_safe` script. Check especially the lines:

  ```
  --------------------------------------------------------------
  MY_PWD=`pwd`
  # Check if we are starting this relative (for the binary release)
  if test -d $MY_PWD/data/mysql -a \
     -f ./share/mysql/english/errmsg.sys -a \
     -x ./bin/mysqld
  --------------------------------------------------------------
  ```

The test performed by these lines should be successful, or you might encounter problems. See Section 4.3.2, “`mysqld_safe` — MySQL Server Startup Script”.

- You might want to use the `--user` option for `mysqld`, but to do this you need to run the `mysqld_multi` script as the Unix superuser (`root`). Having the option in the option file doesn't matter; you just get a warning if you are not the superuser and the `mysqld` processes are started under your own Unix account.

The following example shows how you might set up an option file for use with `mysqld_multi`. The order in which the `mysqld` programs are started or stopped depends on the order in which they appear in the option file. Group numbers need not form an unbroken sequence. The first and fifth `[mysqldN]` groups were intentionally omitted from the example to illustrate that you can have “gaps” in the option file. This gives you more flexibility.

```
# This is an example of a my.cnf file for mysqld_multi.
# Usually this file is located in home dir ~/.my.cnf or /etc/my.cnf

[mysqld_multi]
mysqld     = /usr/local/mysql/bin/mysqld_safe
mysqladmin = /usr/local/mysql/bin/mysqladmin
user       = multi_admin
password   = my_password

[mysqld2]
socket     = /tmp/mysql.sock2
port       = 3307
pid-file   = /usr/local/mysql/data2/hostname.pid2
datadir    = /usr/local/mysql/data2
language   = /usr/local/mysql/share/mysql/english
user       = unix_user1

[mysqld3]
mysqld     = /path/to/mysqld_safe
ledir      = /path/to/mysqld-binary/
mysqladmin = /path/to/mysqladmin
socket     = /tmp/mysql.sock3
port       = 3308
pid-file   = /usr/local/mysql/data3/hostname.pid3
datadir    = /usr/local/mysql/data3
language   = /usr/local/mysql/share/mysql/swedish
user       = unix_user2

[mysqld4]
socket     = /tmp/mysql.sock4
port       = 3309
pid-file   = /usr/local/mysql/data4/hostname.pid4
datadir    = /usr/local/mysql/data4
language   = /usr/local/mysql/share/mysql/estonia
user       = unix_user3

[mysqld6]
socket     = /tmp/mysql.sock6
port       = 3311
pid-file   = /usr/local/mysql/data6/hostname.pid6
datadir    = /usr/local/mysql/data6
language   = /usr/local/mysql/share/mysql/japanese
user       = unix_user4
```

See Section 4.2.6, “Using Option Files”.

4.4 MySQL Installation-Related Programs

The programs in this section are used when installing or upgrading MySQL.

4.4.1 `comp_err` — Compile MySQL Error Message File

comp_err creates the errmsg.sys file that is used by mysqld to determine the error messages to display for different error codes. comp_err normally is run automatically when MySQL is built. It compiles the errmsg.sys file from the text file located at sql/share/errmsg-utf8.txt in MySQL source distributions.

comp_err also generates mysqld_error.h, mysqld_ername.h, and sql_state.h header files.

For more information about how error messages are defined, see the MySQL Internals Manual.

Invoke comp_err like this:

```
shell> comp_err [options]
```

comp_err supports the following options.

- --help, -?

  Display a help message and exit.

- --charset=*dir_name*, -C *dir_name*

  The character set directory. The default is ../sql/share/charsets.

- --debug=*debug_options*, -# *debug_options*

  Write a debugging log. A typical *debug_options* string is d:t:O,*file_name*. The default is d:t:O,/tmp/comp_err.trace.

- --debug-info, -T

  Print some debugging information when the program exits.

- --header_file=*file_name*, -H *file_name*

  The name of the error header file. The default is mysqld_error.h.

- --in_file=*file_name*, -F *file_name*

  The name of the input file. The default is ../sql/share/errmsg-utf8.txt.

- --name_file=*file_name*, -N *file_name*

  The name of the error name file. The default is mysqld_ername.h.

- --out_dir=*dir_name*, -D *dir_name*

  The name of the output base directory. The default is ../sql/share/.

- --out_file=*file_name*, -O *file_name*

  The name of the output file. The default is errmsg.sys.

- --statefile=*file_name*, -S *file_name*

  The name for the SQLSTATE header file. The default is sql_state.h.

- --version, -V

  Display version information and exit.

4.4.2 mysql_secure_installation — Improve MySQL Installation Security

This program enables you to improve the security of your MySQL installation in the following ways:

- You can set a password for root accounts.
- You can remove root accounts that are accessible from outside the local host.
- You can remove anonymous-user accounts.
- You can remove the test database (which by default can be accessed by all users, even anonymous users), and privileges that permit anyone to access databases with names that start with test_.

mysql_secure_installation helps you implement security recommendations similar to those described at Section 2.9.4, “Securing the Initial MySQL Account”.

Normal usage is to connect to the local MySQL server; invoke mysql_secure_installation without arguments:

```
shell> mysql_secure_installation
```

When executed, mysql_secure_installation prompts you to determine which actions to perform.

The validate_password plugin can be used for password strength checking. If the plugin is not installed, mysql_secure_installation prompts the user whether to install it. Any passwords entered later are checked using the plugin if it is enabled.

Most of the usual MySQL client options such as --host and --port can be used on the command line and in option files. For example, to connect to the local server over IPv6 using port 3307, use this command:

```
shell> mysql_secure_installation --host=::1 --port=3307
```

mysql_secure_installation supports the following options, which can be specified on the command line or in the [mysql_secure_installation] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.5 mysql_secure_installation Options

Format	Description
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--help	Display help message and exit
--host	Host to connect to (IP address or host name)
--no-defaults	Read no option files
--password	Accepted but always ignored. Whenever mysql_secure_installation is invoked, the user is prompted for a password, regardless.
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format

Format	Description
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--tls-version	Protocols permitted for secure connections
--use-default	Execute with no user interactivity
--user	MySQL user name to use when connecting to server

- `--help`, `-?`

  Display a help message and exit.

- `--defaults-extra-file=`*`file_name`*

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=`*`file_name`*

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-group-suffix=`*`str`*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysql_secure_installation` normally reads the `[client]` and `[mysql_secure_installation]` groups. If the `--defaults-group-suffix=_other` option is given, `mysql_secure_installation` also reads the `[client_other]` and `[mysql_secure_installation_other]` groups.

- `--host=`*`host_name`*, `-h` *`host_name`*

  Connect to the MySQL server on the given host.

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.)

- `--password=`*`password`*, `-p` *`password`*

  This option is accepted but ignored. Whether or not this option is used, `mysql_secure_installation` always prompts the user for a password.

- `--port=`*`port_num`*, `-P` *`port_num`*

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=protocol_list`

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--use-default`

  Execute noninteractively. This option can be used for unattended installation operations.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server.

4.4.3 mysql_ssl_rsa_setup — Create SSL/RSA Files

This program creates the SSL certificate and key files and RSA key-pair files required to support secure connections using SSL and secure password exchange using RSA over unencrypted connections, if those files are missing. `mysql_ssl_rsa_setup` can also be used to create new SSL files if the existing ones have expired.

Note

`mysql_ssl_rsa_setup` uses the `openssl` command, so its use is contingent on having OpenSSL installed on your machine.

Another way to generate SSL and RSA files, for MySQL distributions compiled using OpenSSL, is to have the server generated them automatically. See Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.

Important

`mysql_ssl_rsa_setup` helps lower the barrier to using SSL by making it easier to generate the required files. However, certificates generated by `mysql_ssl_rsa_setup` are self-signed, which is not very secure. After you gain experience using the files created by `mysql_ssl_rsa_setup`, consider obtaining a CA certificate from a registered certificate authority.

Invoke `mysql_ssl_rsa_setup` like this:

```
shell> mysql_ssl_rsa_setup [options]
```

Typical options are `--datadir` to specify where to create the files, and `--verbose` to see the `openssl` commands that `mysql_ssl_rsa_setup` executes.

mysql_ssl_rsa_setup attempts to create SSL and RSA files using a default set of file names. It works as follows:

1. mysql_ssl_rsa_setup checks for the openssl binary at the locations specified by the PATH environment variable. If openssl is not found, mysql_ssl_rsa_setup does nothing. If openssl is present, mysql_ssl_rsa_setup looks for default SSL and RSA files in the MySQL data directory specified by the --datadir option, or the compiled-in data directory if that option is not given.

2. mysql_ssl_rsa_setup checks the data directory for SSL files with the following names:

```
ca.pem
server-cert.pem
server-key.pem
```

3. If any of those files are present, mysql_ssl_rsa_setup creates no SSL files. Otherwise, it invokes openssl to create them, plus some additional files:

```
ca.pem               Self-signed CA certificate
ca-key.pem           CA private key
server-cert.pem      Server certificate
server-key.pem       Server private key
client-cert.pem      Client certificate
client-key.pem       Client private key
```

   These files enable secure client connections using SSL; see Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

4. mysql_ssl_rsa_setup checks the data directory for RSA files with the following names:

```
private_key.pem      Private member of private/public key pair
public_key.pem       Public member of private/public key pair
```

5. If any of these files are present, mysql_ssl_rsa_setup creates no RSA files. Otherwise, it invokes openssl to create them. These files enable secure password exchange using RSA over unencrypted connections for accounts authenticated by the sha256_password plugin; see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

For information about the characteristics of files created by mysql_ssl_rsa_setup, see Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.

At startup, the MySQL server automatically uses the SSL files created by mysql_ssl_rsa_setup to enable SSL if no explicit SSL options are given other than --ssl. If you prefer to designate the files explicitly, use the --ssl-ca, --ssl-cert, and --ssl-key options at startup to name the ca.pem, server-cert.pem, and server-key.pem files, respectively.

The server also automatically uses the RSA files created by mysql_ssl_rsa_setup to enable RSA if no explicit RSA options are given.

If the server is SSL-enabled, clients use SSL by default for the connection. To specify certificate and key files explicitly, use the --ssl-ca, --ssl-cert, and --ssl-key options to name the ca.pem, client-cert.pem, and client-key.pem files, respectively. However, some additional client setup may be required first because mysql_ssl_rsa_setup by default creates those files in the data directory. The permissions for the data directory normally enable access only to the system account that runs the MySQL server, so client programs cannot use files located there. To make the files available, copy them to a directory that is readable (but *not* writable) by clients:

- For local clients, the MySQL installation directory can be used. For example, if the data directory is a subdirectory of the installation directory and your current location is the data directory, you can copy the files like this:

```
shell> cp ca.pem client-cert.pem client-key.pem ..
```

- For remote clients, distribute the files using a secure channel to ensure they are not tampered with during transit.

If the SSL files used for a MySQL installation have expired, you can use mysql_ssl_rsa_setup to create new ones:

1. Stop the server.
2. Rename or remove the existing SSL files. You may wish to make a backup of them first. (The RSA files do not expire, so you need not remove them. mysql_ssl_rsa_setup will see that they exist and not overwrite them.)
3. Run mysql_ssl_rsa_setup with the --datadir option to specify where to create the new files.
4. Restart the server.

mysql_ssl_rsa_setup supports the following command-line options, which can be specified on the command line or in the [mysql_ssl_rsa_setup] and [mysqld] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, "Using Option Files".

Table 4.6 mysql_ssl_rsa_setup Options

Format	Description
--datadir	Path to data directory
--help	Display help message and exit
--suffix	Suffix for X509 certificate Common Name attribute
--uid	Name of effective user to use for file permissions
--verbose	Verbose mode
--version	Display version information and exit

- --help, ?

  Display a help message and exit.

- --datadir=*dir_name*

  The path to the directory that mysql_ssl_rsa_setup should check for default SSL and RSA files and in which it should create files if they are missing. The default is the compiled-in data directory.

- --suffix=*str*

  The suffix for the Common Name attribute in X509 certificates. The suffix value is limited to 17 characters. The default is based on the MySQL version number.

- --uid=name, -v

  The name of the user who should be the owner of any created files. The value is a user name, not a numeric user ID. In the absence of this option, files created by mysql_ssl_rsa_setup are owned by the user who executes it. This option is valid only if you execute the program as root on a system that supports the chown() system call.

- --verbose, -v

  Verbose mode. Produce more output about what the program does. For example, the program shows the openssl commands it runs, and produces output to indicate whether it skips SSL or RSA file creation because some default file already exists.

- `--version`, `-V`

  Display version information and exit.

4.4.4 mysql_tzinfo_to_sql — Load the Time Zone Tables

The `mysql_tzinfo_to_sql` program loads the time zone tables in the `mysql` database. It is used on systems that have a *zoneinfo* database (the set of files describing time zones). Examples of such systems are Linux, FreeBSD, Solaris, and OS X. One likely location for these files is the `/usr/share/zoneinfo` directory (`/usr/share/lib/zoneinfo` on Solaris). If your system does not have a zoneinfo database, you can use the downloadable package described in Section 10.6, "MySQL Server Time Zone Support".

`mysql_tzinfo_to_sql` can be invoked several ways:

```
shell> mysql_tzinfo_to_sql tz_dir
shell> mysql_tzinfo_to_sql tz_file tz_name
shell> mysql_tzinfo_to_sql --leap tz_file
```

For the first invocation syntax, pass the zoneinfo directory path name to `mysql_tzinfo_to_sql` and send the output into the `mysql` program. For example:

```
shell> mysql_tzinfo_to_sql /usr/share/zoneinfo | mysql -u root mysql
```

`mysql_tzinfo_to_sql` reads your system's time zone files and generates SQL statements from them. `mysql` processes those statements to load the time zone tables.

The second syntax causes `mysql_tzinfo_to_sql` to load a single time zone file *`tz_file`* that corresponds to a time zone name *`tz_name`*:

```
shell> mysql_tzinfo_to_sql tz_file tz_name | mysql -u root mysql
```

If your time zone needs to account for leap seconds, invoke `mysql_tzinfo_to_sql` using the third syntax, which initializes the leap second information. *`tz_file`* is the name of your time zone file:

```
shell> mysql_tzinfo_to_sql --leap tz_file | mysql -u root mysql
```

After running `mysql_tzinfo_to_sql`, it is best to restart the server so that it does not continue to use any previously cached time zone data.

4.4.5 mysql_upgrade — Check and Upgrade MySQL Tables

`mysql_upgrade` examines all tables in all databases for incompatibilities with the current version of MySQL Server. `mysql_upgrade` also upgrades the system tables so that you can take advantage of new privileges or capabilities that might have been added.

If `mysql_upgrade` finds that a table has a possible incompatibility, it performs a table check and, if problems are found, attempts a table repair. If the table cannot be repaired, see Section 2.10.3, "Rebuilding or Repairing Tables or Indexes" for manual table repair strategies.

You should execute `mysql_upgrade` each time you upgrade MySQL.

`mysql_upgrade` communicates directly with the MySQL server, sending it the SQL statements required to perform an upgrade.

Note

On Windows Server 2008, Vista, and newer, you must run `mysql_upgrade` with administrator privileges. You can do this by running a Command Prompt

as Administrator and running the command. Failure to do so may result in the upgrade failing to execute correctly.

Caution

You should always back up your current MySQL installation *before* performing an upgrade. See Section 7.2, “Database Backup Methods”.

Some upgrade incompatibilities may require special handling before you upgrade your MySQL installation and run mysql_upgrade. See Section 2.10.1, “Upgrading MySQL”, for instructions on determining whether any such incompatibilities apply to your installation and how to handle them.

To use mysql_upgrade, make sure that the server is running. Then invoke it like this to check and repair tables and to upgrade the system tables:

```
shell> mysql_upgrade [options]
```

After running mysql_upgrade, stop the server and restart it so that any changes made to the system tables take effect.

If you have multiple MySQL server instances running, invoke mysql_upgrade with connection parameters appropriate for connecting to the desired server. For example, with servers running on the local host on parts 3306 through 3308, upgrade each of them by connecting to the appropriate port:

```
shell> mysql_upgrade --protocol=tcp -P 3306 [other_options]
shell> mysql_upgrade --protocol=tcp -P 3307 [other_options]
shell> mysql_upgrade --protocol=tcp -P 3308 [other_options]
```

For local host connections on Unix, the --protocol=tcp option forces a connection using TCP/IP rather than the Unix socket file.

mysql_upgrade processes all tables in all databases, which might take a long time to complete. Each table is locked and therefore unavailable to other sessions while it is being processed. Check and repair operations can be time-consuming, particularly for large tables.

For details about what table-checking operations entail, see the description of the FOR UPGRADE option of the CHECK TABLE statement (see Section 13.7.2.2, “CHECK TABLE Syntax”).

All checked and repaired tables are marked with the current MySQL version number. This ensures that next time you run mysql_upgrade with the same version of the server, it can tell whether there is any need to check or repair the table again.

mysql_upgrade also saves the MySQL version number in a file named mysql_upgrade_info in the data directory. This is used to quickly check whether all tables have been checked for this release so that table-checking can be skipped. To ignore this file and perform the check regardless, use the --force option.

mysql_upgrade checks user table rows and, for any row with an empty plugin column, sets that column to 'mysql_native_password' if the credentials use a hash format compatible with that plugin. Rows with a pre-4.1 password hash must be upgraded manually.

mysql_upgrade does not upgrade the contents of the help tables. For upgrade instructions, see Section 5.1.10, “Server-Side Help”.

Unless invoked with the --skip-sys-schema option, mysql_upgrade installs the sys schema if it is not installed, and upgrades it to the current version otherwise. mysql_upgrade returns an error if a sys schema exists but has no version view, on the assumption that its absence indicates a user-created schema:

```
Error occurred: A sys schema exists with no sys.version view. If
```

```
you have a user created sys schema, this must be renamed for the
upgrade to succeed.
```

To upgrade in this case, remove or rename the existing sys schema first.

mysql_upgrade checks for partitioned InnoDB tables that were created using the generic partitioning handler and attempts to upgrade them to InnoDB native partitioning. You can upgrade such tables individually in the mysql client using the ALTER TABLE ... UPGRADE PARTITIONING SQL statement.

By default, mysql_upgrade runs as the MySQL root user. If the root password is expired when you run mysql_upgrade, you will see a message that your password is expired and that mysql_upgrade failed as a result. To correct this, reset the root password to unexpire it and run mysql_upgrade again. First, connect to the server as root:

```
shell> mysql -u root -p
Enter password: ****  <- enter root password here
```

Reset the password using ALTER USER:

```
mysql> ALTER USER USER() IDENTIFIED BY 'root-password';
```

Then exit mysql and run mysql_upgrade again:

```
shell> mysql_upgrade [options]
```

mysql_upgrade supports the following options, which can be specified on the command line or in the [mysql_upgrade] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.7 mysql_upgrade Options

Format	Description
--bind-address	Use specified network interface to connect to MySQL Server
--character-sets-dir	Directory where character sets are installed
--compress	Compress all information sent between client and server
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--default-character-set	Specify default character set
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--force	Force execution even if mysql_upgrade has already been executed for current version of MySQL
--help	Display help message and exit
--host	Connect to MySQL server on given host
--login-path	Read login path options from .mylogin.cnf
--max-allowed-packet	Maximum packet length to send to or receive from server
--net-buffer-length	Buffer size for TCP/IP and socket communication
--no-defaults	Read no option files

Format	Description
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--skip-sys-schema	Do not install or upgrade the sys schema
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tls-version	Protocols permitted for secure connections
--upgrade-system-tables	Update only system tables, not data
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version-check	Check for proper server version
--write-binlog	Write all statements to binary log

- `--help`

  Display a short help message and exit.

- `--basedir=`*`dir_name`*

  The path to the MySQL installation directory.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--debug[=`*`debug_options`*`]`, `-# [`*`debug_options`*`]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:O,/tmp/mysql_upgrade.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`, `-T`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, “Character Set Configuration”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysql_upgrade` normally reads the `[client]` and `[mysql_upgrade]` groups. If the `--defaults-group-suffix=_other` option is given, `mysql_upgrade` also reads the `[client_other]` and `[mysql_upgrade_other]` groups.

- `--force`

  Ignore the `mysql_upgrade_info` file and force execution even if `mysql_upgrade` has already been executed for the current version of MySQL.

- `--host=host_name`, `-h host_name`

  Connect to the MySQL server on the given host.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--max-allowed-packet=value`

  The maximum size of the buffer for client/server communication. The default value is 24MB. The minimum and maximum values are 4KB and 2GB.

- `--net-buffer-length=value`

  The initial size of the buffer for client/server communication. The default value is 1MB − 1KB. The minimum and maximum values are 4KB and 16MB.

- `--no-defaults`

Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, "`mysql_config_editor` — MySQL Configuration Utility".)

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysql_upgrade` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, "End-User Guidelines for Password Security". You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysql_upgrade` does not find it. See Section 6.3.10, "Pluggable Authentication".

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, "Connecting to the MySQL Server".

- `--shared-memory-base-name=name`

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--skip-sys-schema`

  `mysql_upgrade` installs the `sys` schema if it is not installed, and upgrades it to the current version otherwise. The `--skip-sys-schema` option suppresses this behavior.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=protocol_list`

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--upgrade-system-tables`, `-s`

  Upgrade only the system tables, do not upgrade data.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server. The default user name is `root`.

- `--verbose`

  Verbose mode. Print more information about what the program does.

- `--version-check`, `-k`

  Check the version of the server to which `mysql_upgrade` is connecting to verify that it is the same as the version for which `mysql_upgrade` was built. If not, `mysql_upgrade` exits. This option is enabled by default; to disable the check, use `--skip-version-check`.

- `--write-binlog`

  By default, binary logging by `mysql_upgrade` is disabled. Invoke the program with `--write-binlog` if you want its actions to be written to the binary log.

Note

this option is not recommended on MySQL Servers running with `--gtid-mode=ON` because `mysql_upgrade` can make changes to system tables that use the `MyISAM` storage engine, which is non-transactional.

4.5 MySQL Client Programs

This section describes client programs that connect to the MySQL server.

4.5.1 `mysql` — The MySQL Command-Line Tool

`mysql` is a simple SQL shell with input line editing capabilities. It supports interactive and noninteractive use. When used interactively, query results are presented in an ASCII-table format. When used noninteractively (for example, as a filter), the result is presented in tab-separated format. The output format can be changed using command options.

If you have problems due to insufficient memory for large result sets, use the `--quick` option. This forces `mysql` to retrieve results from the server a row at a time rather than retrieving the entire result set and buffering it in memory before displaying it. This is done by returning the result set using the `mysql_use_result()` C API function in the client/server library rather than `mysql_store_result()`.

Note

Alternatively, MySQL Shell offers access to the X DevAPI. For details, see Chapter 19, *MySQL Shell User Guide*.

Using mysql is very easy. Invoke it from the prompt of your command interpreter as follows:

```
shell> mysql db_name
```

Or:

```
shell> mysql --user=user_name --password db_name
Enter password: your_password
```

Then type an SQL statement, end it with ;, \g, or \G and press Enter.

Typing **Control+C** interrupts the current statement if there is one, or cancels any partial input line otherwise.

You can execute SQL statements in a script file (batch file) like this:

```
shell> mysql db_name < script.sql > output.tab
```

On Unix, the mysql client logs statements executed interactively to a history file. See Section 4.5.1.3, “mysql Logging”.

4.5.1.1 mysql Options

mysql supports the following options, which can be specified on the command line or in the [mysql] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.8 mysql Options

Format	Description	Introduced
--auto-rehash	Enable automatic rehashing	
--auto-vertical-output	Enable automatic vertical result set display	
--batch	Do not use history file	
--binary-as-hex	Display binary values in hexadecimal notation	8.0.2
--binary-mode	Disable \r\n - to - \n translation and treatment of \0 as end-of-query	
--bind-address	Use specified network interface to connect to MySQL Server	
--character-sets-dir	Directory where character sets are installed	
--column-names	Write column names in results	
--column-type-info	Display result set metadata	
--comments	Whether to retain or strip comments in statements sent to the server	
--compress	Compress all information sent between client and server	
--connect-expired-password	Indicate to server that client can handle expired-password sandbox mode.	
--connect_timeout	Number of seconds before connection timeout	
--database	The database to use	
--debug	Write debugging log; supported only if MySQL was built with debugging support	
--debug-check	Print debugging information when program exits	
--debug-info	Print debugging information, memory, and CPU statistics when program exits	

Format	Description	Introduced
--default-auth	Authentication plugin to use	
--default-character-set	Specify default character set	
--defaults-extra-file	Read named option file in addition to usual option files	
--defaults-file	Read only named option file	
--defaults-group-suffix	Option group suffix value	
--delimiter	Set the statement delimiter	
--enable-cleartext-plugin	Enable cleartext authentication plugin	
--execute	Execute the statement and quit	
--force	Continue even if an SQL error occurs	
--help	Display help message and exit	
--histignore	Patterns specifying which statements to ignore for logging	
--host	Connect to MySQL server on given host	
--html	Produce HTML output	
--ignore-spaces	Ignore spaces after function names	
--init-command	SQL statement to execute after connecting	
--line-numbers	Write line numbers for errors	
--local-infile	Enable or disable for LOCAL capability for LOAD DATA INFILE	
--login-path	Read login path options from .mylogin.cnf	
--max_allowed_packet	Maximum packet length to send to or receive from server	
--max_join_size	The automatic limit for rows in a join when using --safe-updates	
--named-commands	Enable named mysql commands	
--net_buffer_length	Buffer size for TCP/IP and socket communication	
--no-auto-rehash	Disable automatic rehashing	
--no-beep	Do not beep when errors occur	
--no-defaults	Read no option files	
--one-database	Ignore statements except those for the default database named on the command line	
--pager	Use the given command for paging query output	
--password	Password to use when connecting to server	
--pipe	On Windows, connect to server using named pipe	
--plugin-dir	Directory where plugins are installed	
--port	TCP/IP port number to use for connection	
--print-defaults	Print default options	
--prompt	Set the prompt to the specified format	
--protocol	Connection protocol to use	
--quick	Do not cache each query result	
--raw	Write column values without escape conversion	
--reconnect	If the connection to the server is lost, automatically try to reconnect	

Format	Description	Introduced
--i-am-a-dummy, --safe-updates	Allow only UPDATE and DELETE statements that specify key values	
--secure-auth	Do not send passwords to server in old (pre-4.1) format	
--select_limit	The automatic limit for SELECT statements when using --safe-updates	
--server-public-key-path	Path name to file containing RSA public key	
--shared-memory-base-name	The name of shared memory to use for shared-memory connections	
--show-warnings	Show warnings after each statement if there are any	
--sigint-ignore	Ignore SIGINT signals (typically the result of typing Control+C)	
--silent	Silent mode	
--skip-auto-rehash	Disable automatic rehashing	
--skip-column-names	Do not write column names in results	
--skip-line-numbers	Skip line numbers for errors	
--skip-named-commands	Disable named mysql commands	
--skip-pager	Disable paging	
--skip-reconnect	Disable reconnecting	
--socket	For connections to localhost, the Unix socket file or Windows named pipe to use	
--ssl-ca	Path of file that contains list of trusted SSL CAs	
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format	
--ssl-cert	Path of file that contains X509 certificate in PEM format	
--ssl-cipher	List of permitted ciphers to use for connection encryption	
--ssl-crl	Path of file that contains certificate revocation lists	
--ssl-crlpath	Path of directory that contains certificate revocation list files	
--ssl-key	Path of file that contains X509 key in PEM format	
--ssl-mode	Security state of connection to server	
--syslog	Log interactive statements to syslog	
--table	Display output in tabular format	
--tee	Append a copy of output to named file	
--tls-version	Protocols permitted for secure connections	
--unbuffered	Flush the buffer after each query	
--user	MySQL user name to use when connecting to server	
--verbose	Verbose mode	
--version	Display version information and exit	
--vertical	Print query output rows vertically (one line per column value)	
--wait	If the connection cannot be established, wait and retry instead of aborting	
--xml	Produce XML output	

- `--help`, `-?`

  Display a help message and exit.

- `--auto-rehash`

  Enable automatic rehashing. This option is on by default, which enables database, table, and column name completion. Use `--disable-auto-rehash` to disable rehashing. That causes `mysql` to start faster, but you must issue the `rehash` command or its `\#` shortcut if you want to use name completion.

  To complete a name, enter the first part and press Tab. If the name is unambiguous, `mysql` completes it. Otherwise, you can press Tab again to see the possible names that begin with what you have typed so far. Completion does not occur if there is no default database.

  Note

  This feature requires a MySQL client that is compiled with the **readline** library. Typically, the **readline** library is not available on Windows.

- `--auto-vertical-output`

  Cause result sets to be displayed vertically if they are too wide for the current window, and using normal tabular format otherwise. (This applies to statements terminated by `;` or `\G`.)

- `--batch`, `-B`

  Print results using tab as the column separator, with each row on a new line. With this option, `mysql` does not use the history file.

  Batch mode results in nontabular output format and escaping of special characters. Escaping may be disabled by using raw mode; see the description for the `--raw` option.

- `--binary-as-hex`, `-b`

  When this option is given, `mysql` displays binary data using hexadecimal notation (`0x`*`value`*). This occurs whether the overall output dislay format is tabular, vertical, HTML, or XML.

  This option was added in MySQL 8.0.2.

- `--binary-mode`

  This option helps when processing `mysqlbinlog` output that may contain `BLOB` values. By default, `mysql` translates `\r\n` in statement strings to `\n` and interprets `\0` as the statement terminator. `--binary-mode` disables both features. It also disables all `mysql` commands except `charset` and `delimiter` in non-interactive mode (for input piped to `mysql` or loaded using the `source` command).

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--column-names`

  Write column names in results.

- `--column-type-info`

Display result set metadata.

- `--comments`, `-c`

  Whether to strip or preserve comments in statements sent to the server. The default is `--skip-comments` (strip comments), enable with `--comments` (preserve comments).

  Note

  The `mysql` client always passes optimizer hints to the server, regardless of whether this option is given.

  Comment stripping is deprecated. This feature and the options to control it will be removed in a future MySQL release.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--connect-expired-password`

  Indicate to the server that the client can handle sandbox mode if the account used to connect has an expired password. This can be useful for noninteractive invocations of `mysql` because normally the server disconnects noninteractive clients that attempt to connect using an account with an expired password. (See Section 6.3.9, “Password Expiration and Sandbox Mode”.)

- `--database=`*`db_name`*, `-D` *`db_name`*

  The database to use. This is useful primarily in an option file.

- `--debug[=`*`debug_options`*`]`, `-# [`*`debug_options`*`]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:o,/tmp/mysql.trace`.

  This option is available only if MySQL was built using `WITH_DEBUG`. MySQL release binaries provided by Oracle are *not* built using this option.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`, `-T`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=`*`plugin`*

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--default-character-set=`*`charset_name`*

  Use *`charset_name`* as the default character set for the client and connection.

  This option can be useful if the operating system uses one character set and the `mysql` client by default uses another. In this case, output may be formatted incorrectly. You can usually fix such issues by using this option to force the client to use the system character set instead.

  For more information, see Section 10.1.4, “Connection Character Sets and Collations”, and Section 10.5, “Character Set Configuration”.

- `--defaults-extra-file=`*`file_name`*

Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *file_name* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- --defaults-file=*file_name*

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *file_name* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with --defaults-file, client programs read .mylogin.cnf.

- --defaults-group-suffix=*str*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *str*. For example, mysql normally reads the [client] and [mysql] groups. If the --defaults-group-suffix=_other option is given, mysql also reads the [client_other] and [mysql_other] groups.

- --delimiter=*str*

  Set the statement delimiter. The default is the semicolon character (;).

- --disable-named-commands

  Disable named commands. Use the * form only, or use named commands only at the beginning of a line ending with a semicolon (;). mysql starts with this option *enabled* by default. However, even with this option, long-format commands still work from the first line. See Section 4.5.1.2, “mysql Commands”.

- --enable-cleartext-plugin

  Enable the mysql_clear_password cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- --execute=*statement*, -e *statement*

  Execute the statement and quit. The default output format is like that produced with --batch. See Section 4.2.4, “Using Options on the Command Line”, for some examples. With this option, mysql does not use the history file.

- --force, -f

  Continue even if an SQL error occurs.

- --histignore

  A colon-separated list of one or more patterns specifying statements to ignore for logging purposes. These patterns are added to the default pattern list ("*IDENTIFIED*:*PASSWORD*"). The value specified for this option affects logging of statements written to the history file, and to syslog if the --syslog option is given. For more information, see Section 4.5.1.3, “mysql Logging”.

- --host=*host_name*, -h *host_name*

  Connect to the MySQL server on the given host.

- --html, -H

  Produce HTML output.

- --ignore-spaces, -i

Ignore spaces after function names. The effect of this is described in the discussion for the IGNORE_SPACE SQL mode (see Section 5.1.8, “Server SQL Modes”).

- --init-command=str

  SQL statement to execute after connecting to the server. If auto-reconnect is enabled, the statement is executed again after reconnection occurs.

- --line-numbers

  Write line numbers for errors. Disable this with --skip-line-numbers.

- --local-infile[={0|1}]

  Enable or disable LOCAL capability for LOAD DATA INFILE. For mysql, this capability is disabled by default. With no value, the option enables LOCAL. The option may be given as --local-infile=0 or --local-infile=1 to explicitly disable or enable LOCAL. Enabling local data loading also requires that the server permits it; see Section 6.1.6, “Security Issues with LOAD DATA LOCAL”

- --login-path=*name*

  Read options from the named login path in the .mylogin.cnf login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

- --named-commands, -G

  Enable named mysql commands. Long-format commands are permitted, not just short-format commands. For example, quit and \q both are recognized. Use --skip-named-commands to disable named commands. See Section 4.5.1.2, “mysql Commands”.

- --no-auto-rehash, -A

  This has the same effect as --skip-auto-rehash. See the description for --auto-rehash.

- --no-beep, -b

  Do not beep when errors occur.

- --no-defaults

  Do not read any option files. If program startup fails due to reading unknown options from an option file, --no-defaults can be used to prevent them from being read.

  The exception is that the .mylogin.cnf file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when --no-defaults is used. (.mylogin.cnf is created by the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.)

- --one-database, -o

  Ignore statements except those that occur while the default database is the one named on the command line. This option is rudimentary and should be used with care. Statement filtering is based only on USE statements.

  Initially, mysql executes statements in the input because specifying a database *db_name* on the command line is equivalent to inserting USE *db_name* at the beginning of the input. Then, for each USE statement encountered, mysql accepts or rejects following statements depending on whether the database named is the one on the command line. The content of the statements is immaterial.

  Suppose that mysql is invoked to process this set of statements:

```
DELETE FROM db2.t2;
USE db2;
DROP TABLE db1.t1;
CREATE TABLE db1.t1 (i INT);
USE db1;
INSERT INTO t1 (i) VALUES(1);
CREATE TABLE db2.t1 (j INT);
```

If the command line is `mysql --force --one-database db1`, `mysql` handles the input as follows:

- The `DELETE` statement is executed because the default database is `db1`, even though the statement names a table in a different database.
- The `DROP TABLE` and `CREATE TABLE` statements are not executed because the default database is not `db1`, even though the statements name a table in `db1`.
- The `INSERT` and `CREATE TABLE` statements are executed because the default database is `db1`, even though the `CREATE TABLE` statement names a table in a different database.

- `--pager[=command]`

  Use the given command for paging query output. If the command is omitted, the default pager is the value of your `PAGER` environment variable. Valid pagers are `less`, `more`, `cat [> filename]`, and so forth. This option works only on Unix and only in interactive mode. To disable paging, use `--skip-pager`. Section 4.5.1.2, “mysql Commands”, discusses output paging further.

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysql` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysql` does not find it. See Section 6.3.10, “Pluggable Authentication”.

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--prompt=format_str`

  Set the prompt to the specified format. The default is `mysql>`. The special sequences that the prompt can contain are described in Section 4.5.1.2, “mysql Commands”.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--quick`, `-q`

  Do not cache each query result, print each row as it is received. This may slow down the server if the output is suspended. With this option, `mysql` does not use the history file.

- `--raw`, `-r`

  For tabular output, the “boxing” around columns enables one column value to be distinguished from another. For nontabular output (such as is produced in batch mode or when the `--batch` or `--silent` option is given), special characters are escaped in the output so they can be identified easily. Newline, tab, `NUL`, and backslash are written as `\n`, `\t`, `\0`, and `\\`. The `--raw` option disables this character escaping.

  The following example demonstrates tabular versus nontabular output and the use of raw mode to disable escaping:

```
% mysql
mysql> SELECT CHAR(92);
+----------+
| CHAR(92) |
+----------+
| \        |
+----------+

% mysql -s
mysql> SELECT CHAR(92);
CHAR(92)
\\

% mysql -s -r
mysql> SELECT CHAR(92);
CHAR(92)
\
```

- `--reconnect`

  If the connection to the server is lost, automatically try to reconnect. A single reconnect attempt is made each time the connection is lost. To suppress reconnection behavior, use `--skip-reconnect`.

- `--safe-updates`, `--i-am-a-dummy`, `-U`

  Permit only those `UPDATE` and `DELETE` statements that specify which rows to modify by using key values. If you have set this option in an option file, you can override it by using `--safe-updates` on the command line. See Section 4.5.1.6, “mysql Tips”, for more information about this option.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

Note

Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be

avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--server-public-key-path=file_name`

  The path name to a file containing the server RSA public key. The file must be in PEM format. The public key is used for RSA encryption of the client password for connections to the server made using accounts that authenticate with the `sha256_password` plugin. This option is ignored for client accounts that do not authenticate with that plugin. It is also ignored if password encryption is not needed, as is the case when the client connects to the server using an SSL connection.

  The server sends the public key to the client as needed, so it is not necessary to use this option for RSA password encryption to occur. It is more efficient to do so because then the server need not send the key.

  For additional discussion regarding use of the `sha256_password` plugin, including how to get the RSA public key, see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

  This option is available only if MySQL was built using OpenSSL.

- `--shared-memory-base-name=`*`name`*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--show-warnings`

  Cause warnings to be shown after each statement if there are any. This option applies to interactive and batch mode.

- `--sigint-ignore`

  Ignore `SIGINT` signals (typically the result of typing **Control+C**).

- `--silent`, `-s`

  Silent mode. Produce less output. This option can be given multiple times to produce less and less output.

  This option results in nontabular output format and escaping of special characters. Escaping may be disabled by using raw mode; see the description for the `--raw` option.

- `--skip-column-names`, `-N`

  Do not write column names in results.

- `--skip-line-numbers`, `-L`

  Do not write line numbers for errors. Useful when you want to compare result files that include error messages.

- `--socket=`*`path`*, `-S` *`path`*

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--syslog`, `-j`

  This option causes `mysql` to send interactive statements to the system logging facility. On Unix, this is `syslog`; on Windows, it is the Windows Event Log. The destination where logged messages appear is system dependent. On Linux, the destination is often the `/var/log/messages` file.

  Here is a sample of output generated on Linux by using `--syslog`. This output is formatted for readability; each logged message actually takes a single line.

```
Mar  7 12:39:25 myhost MysqlClient[20824]:
  SYSTEM_USER:'oscar', MYSQL_USER:'my_oscar', CONNECTION_ID:23,
  DB_SERVER:'127.0.0.1', DB:'--', QUERY:'USE test;'
Mar  7 12:39:28 myhost MysqlClient[20824]:
  SYSTEM_USER:'oscar', MYSQL_USER:'my_oscar', CONNECTION_ID:23,
  DB_SERVER:'127.0.0.1', DB:'test', QUERY:'SHOW TABLES;'
```

  For more information, see Section 4.5.1.3, “mysql Logging”.

- `--table`, `-t`

  Display output in table format. This is the default for interactive use, but can be used to produce table output in batch mode.

- `--tee=`*`file_name`*

  Append a copy of output to the given file. This option works only in interactive mode. Section 4.5.1.2, “mysql Commands”, discusses tee files further.

- `--tls-version=`*`protocol_list`*

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--unbuffered`, `-n`

  Flush the buffer after each query.

- `--user=`*`user_name`*, `-u` *`user_name`*

  The MySQL user name to use when connecting to the server.

- `--verbose`, `-v`

  Verbose mode. Produce more output about what the program does. This option can be given multiple times to produce more and more output. (For example, `-v -v -v` produces table output format even in batch mode.)

- `--version`, `-V`

  Display version information and exit.

- `--vertical`, `-E`

  Print query output rows vertically (one line per column value). Without this option, you can specify vertical output for individual statements by terminating them with `\G`.

- `--wait`, `-w`

If the connection cannot be established, wait and retry instead of aborting.

- `--xml`, `-X`

  Produce XML output.

```
<field name="column_name">NULL</field>
```

  The output when `--xml` is used with `mysql` matches that of `mysqldump --xml`. See Section 4.5.4, “mysqldump — A Database Backup Program”, for details.

  The XML output also uses an XML namespace, as shown here:

```
shell> mysql --xml -uroot -e "SHOW VARIABLES LIKE 'version%'"
<?xml version="1.0"?>

<resultset statement="SHOW VARIABLES LIKE 'version%'" xmlns:xsi="http://www.w3.org/2001/XMLSchema-ins
<row>
<field name="Variable_name">version</field>
<field name="Value">5.0.40-debug</field>
</row>

<row>
<field name="Variable_name">version_comment</field>
<field name="Value">Source distribution</field>
</row>

<row>
<field name="Variable_name">version_compile_machine</field>
<field name="Value">i686</field>
</row>

<row>
<field name="Variable_name">version_compile_os</field>
<field name="Value">suse-linux-gnu</field>
</row>
</resultset>
```

  (See Bug #25946.)

You can also set the following variables by using `--var_name=value`.

- `connect_timeout`

  The number of seconds before connection timeout. (Default value is `0`.)

- `max_allowed_packet`

  The maximum size of the buffer for client/server communication. The default is 16MB, the maximum is 1GB.

- `max_join_size`

  The automatic limit for rows in a join when using `--safe-updates`. (Default value is 1,000,000.)

- `net_buffer_length`

  The buffer size for TCP/IP and socket communication. (Default value is 16KB.)

- `select_limit`

  The automatic limit for `SELECT` statements when using `--safe-updates`. (Default value is 1,000.)

4.5.1.2 mysql Commands

mysql sends each SQL statement that you issue to the server to be executed. There is also a set of commands that mysql itself interprets. For a list of these commands, type help or \h at the mysql> prompt:

```
mysql> help

List of all MySQL commands:
Note that all text commands must be first on line and end with ';'
?         (\?) Synonym for `help'.
clear     (\c) Clear the current input statement.
connect   (\r) Reconnect to the server. Optional arguments are db and host.
delimiter (\d) Set statement delimiter.
edit      (\e) Edit command with $EDITOR.
ego       (\G) Send command to mysql server, display result vertically.
exit      (\q) Exit mysql. Same as quit.
go        (\g) Send command to mysql server.
help      (\h) Display this help.
nopager   (\n) Disable pager, print to stdout.
notee     (\t) Don't write into outfile.
pager     (\P) Set PAGER [to_pager]. Print the query results via PAGER.
print     (\p) Print current command.
prompt    (\R) Change your mysql prompt.
quit      (\q) Quit mysql.
rehash    (\#) Rebuild completion hash.
source    (\.) Execute an SQL script file. Takes a file name as an argument.
status    (\s) Get status information from the server.
system    (\!) Execute a system shell command.
tee       (\T) Set outfile [to_outfile]. Append everything into given
               outfile.
use       (\u) Use another database. Takes database name as argument.
charset   (\C) Switch to another charset. Might be needed for processing
               binlog with multi-byte charsets.
warnings  (\W) Show warnings after every statement.
nowarning (\w) Don't show warnings after every statement.
resetconnection(\x) Clean session context.

For server side help, type 'help contents'
```

If mysql is invoked with the --binary-mode option, all mysql commands are disabled except charset and delimiter in non-interactive mode (for input piped to mysql or loaded using the source command).

Each command has both a long and short form. The long form is not case sensitive; the short form is. The long form can be followed by an optional semicolon terminator, but the short form should not.

The use of short-form commands within multiple-line /* ... */ comments is not supported.

- help [*arg*], \h [*arg*], \? [*arg*], ? [*arg*]

  Display a help message listing the available mysql commands.

  If you provide an argument to the help command, mysql uses it as a search string to access server-side help from the contents of the MySQL Reference Manual. For more information, see Section 4.5.1.4, “mysql Server-Side Help”.

- charset *charset_name*, \C *charset_name*

  Change the default character set and issue a SET NAMES statement. This enables the character set to remain synchronized on the client and server if mysql is run with auto-reconnect enabled (which is not recommended), because the specified character set is used for reconnects.

- clear, \c

  Clear the current input. Use this if you change your mind about executing the statement that you are entering.

- `connect [db_name host_name]`, `\r [db_name host_name]`

  Reconnect to the server. The optional database name and host name arguments may be given to specify the default database or the host where the server is running. If omitted, the current values are used.

- `delimiter str`, `\d str`

  Change the string that `mysql` interprets as the separator between SQL statements. The default is the semicolon character (`;`).

  The delimiter string can be specified as an unquoted or quoted argument on the `delimiter` command line. Quoting can be done with either single quote (`'`), double quote (`"`), or backtick (`` ` ``) characters. To include a quote within a quoted string, either quote the string with a different quote character or escape the quote with a backslash (`\`) character. Backslash should be avoided outside of quoted strings because it is the escape character for MySQL. For an unquoted argument, the delimiter is read up to the first space or end of line. For a quoted argument, the delimiter is read up to the matching quote on the line.

  `mysql` interprets instances of the delimiter string as a statement delimiter anywhere it occurs, except within quoted strings. Be careful about defining a delimiter that might occur within other words. For example, if you define the delimiter as `X`, you will be unable to use the word `INDEX` in statements. `mysql` interprets this as `INDE` followed by the delimiter `X`.

  When the delimiter recognized by `mysql` is set to something other than the default of `;`, instances of that character are sent to the server without interpretation. However, the server itself still interprets `;` as a statement delimiter and processes statements accordingly. This behavior on the server side comes into play for multiple-statement execution (see Section 27.7.20, “C API Support for Multiple Statement Execution”), and for parsing the body of stored procedures and functions, triggers, and events (see Section 23.1, “Defining Stored Programs”).

- `edit`, `\e`

  Edit the current input statement. `mysql` checks the values of the `EDITOR` and `VISUAL` environment variables to determine which editor to use. The default editor is `vi` if neither variable is set.

  The `edit` command works only in Unix.

- `ego`, `\G`

  Send the current statement to the server to be executed and display the result using vertical format.

- `exit`, `\q`

  Exit `mysql`.

- `go`, `\g`

  Send the current statement to the server to be executed.

- `nopager`, `\n`

  Disable output paging. See the description for `pager`.

  The `nopager` command works only in Unix.

- `notee`, `\t`

  Disable output copying to the tee file. See the description for `tee`.

- `nowarning`, `\w`

  Disable display of warnings after each statement.

- `pager [command]`, `\P [command]`

  Enable output paging. By using the `--pager` option when you invoke `mysql`, it is possible to browse or search query results in interactive mode with Unix programs such as `less`, `more`, or any other similar program. If you specify no value for the option, `mysql` checks the value of the `PAGER` environment variable and sets the pager to that. Pager functionality works only in interactive mode.

  Output paging can be enabled interactively with the `pager` command and disabled with `nopager`. The command takes an optional argument; if given, the paging program is set to that. With no argument, the pager is set to the pager that was set on the command line, or `stdout` if no pager was specified.

  Output paging works only in Unix because it uses the `popen()` function, which does not exist on Windows. For Windows, the `tee` option can be used instead to save query output, although it is not as convenient as `pager` for browsing output in some situations.

- `print`, `\p`

  Print the current input statement without executing it.

- `prompt [str]`, `\R [str]`

  Reconfigure the `mysql` prompt to the given string. The special character sequences that can be used in the prompt are described later in this section.

  If you specify the `prompt` command with no argument, `mysql` resets the prompt to the default of `mysql>`.

- `quit`, `\q`

  Exit `mysql`.

- `rehash`, `\#`

  Rebuild the completion hash that enables database, table, and column name completion while you are entering statements. (See the description for the `--auto-rehash` option.)

- `resetconnection`, `\x`

  Reset the connection to clear the session state.

  Resetting a connection has effects similar to `mysql_change_user()` or an auto-reconnect except that the connection is not closed and reopened, and re-authentication is not done. See Section 27.7.7.3, “mysql_change_user()”) and see Section 27.7.19, “Controlling Automatic Reconnection Behavior”).

  This example shows how `resetconnection` clears a value maintained in the session state:

```
mysql> SELECT LAST_INSERT_ID(3);
+------------------+
| LAST_INSERT_ID(3) |
+------------------+
|                3 |
+------------------+

mysql> SELECT LAST_INSERT_ID();
+-----------------+
| LAST_INSERT_ID() |
+-----------------+
|               3 |
+-----------------+

mysql> resetconnection;
```

```
mysql> SELECT LAST_INSERT_ID();
+------------------+
| LAST_INSERT_ID() |
+------------------+
|                0 |
+------------------+
```

- `source file_name`, `\. file_name`

  Read the named file and executes the statements contained therein. On Windows, you can specify path name separators as `/` or `\\`.

- `status`, `\s`

  Provide status information about the connection and the server you are using. If you are running in `--safe-updates` mode, `status` also prints the values for the `mysql` variables that affect your queries.

- `system command`, `\! command`

  Execute the given command using your default command interpreter.

  The `system` command works only in Unix.

- `tee [file_name]`, `\T [file_name]`

  By using the `--tee` option when you invoke `mysql`, you can log statements and their output. All the data displayed on the screen is appended into a given file. This can be very useful for debugging purposes also. `mysql` flushes results to the file after each statement, just before it prints its next prompt. Tee functionality works only in interactive mode.

  You can enable this feature interactively with the `tee` command. Without a parameter, the previous file is used. The `tee` file can be disabled with the `notee` command. Executing `tee` again re-enables logging.

- `use db_name`, `\u db_name`

  Use `db_name` as the default database.

- `warnings`, `\W`

  Enable display of warnings after each statement (if there are any).

Here are a few tips about the `pager` command:

- You can use it to write to a file and the results go only to the file:

```
mysql> pager cat > /tmp/log.txt
```

  You can also pass any options for the program that you want to use as your pager:

```
mysql> pager less -n -i -S
```

- In the preceding example, note the `-S` option. You may find it very useful for browsing wide query results. Sometimes a very wide result set is difficult to read on the screen. The `-S` option to `less` can make the result set much more readable because you can scroll it horizontally using the left-arrow and right-arrow keys. You can also use `-S` interactively within `less` to switch the horizontal-browse mode on and off. For more information, read the `less` manual page:

```
shell> man less
```

- The `-F` and `-X` options may be used with `less` to cause it to exit if output fits on one screen, which is convenient when no scrolling is necessary:

```
mysql> pager less -n -i -S -F -X
```

- You can specify very complex pager commands for handling query output:

```
mysql> pager cat | tee /dr1/tmp/res.txt \
          | tee /dr2/tmp/res2.txt | less -n -i -S
```

  In this example, the command would send query results to two files in two different directories on two different file systems mounted on `/dr1` and `/dr2`, yet still display the results onscreen using `less`.

You can also combine the `tee` and `pager` functions. Have a `tee` file enabled and `pager` set to `less`, and you are able to browse the results using the `less` program and still have everything appended into a file the same time. The difference between the Unix `tee` used with the `pager` command and the `mysql` built-in `tee` command is that the built-in `tee` works even if you do not have the Unix `tee` available. The built-in `tee` also logs everything that is printed on the screen, whereas the Unix `tee` used with `pager` does not log quite that much. Additionally, `tee` file logging can be turned on and off interactively from within `mysql`. This is useful when you want to log some queries to a file, but not others.

The `prompt` command reconfigures the default `mysql>` prompt. The string for defining the prompt can contain the following special sequences.

Option	Description
`\C`	The current connection identifier
`\c`	A counter that increments for each statement you issue
`\D`	The full current date
`\d`	The default database
`\h`	The server host
`\l`	The current delimiter
`\m`	Minutes of the current time
`\n`	A newline character
`\O`	The current month in three-letter format (Jan, Feb, …)
`\o`	The current month in numeric format
`\P`	am/pm
`\p`	The current TCP/IP port or socket file
`\R`	The current time, in 24-hour military time (0–23)
`\r`	The current time, standard 12-hour time (1–12)
`\S`	Semicolon
`\s`	Seconds of the current time
`\t`	A tab character
`\U`	Your full *`user_name@host_name`* account name
`\u`	Your user name
`\v`	The server version
`\w`	The current day of the week in three-letter format (Mon, Tue, …)
`\Y`	The current year, four digits
`\y`	The current year, two digits

Option	Description
_	A space
\	A space (a space follows the backslash)
\'	Single quote
\"	Double quote
\\	A literal \ backslash character
\x	x, for any "x" not listed above

You can set the prompt in several ways:

- *Use an environment variable.* You can set the MYSQL_PS1 environment variable to a prompt string. For example:

```
shell> export MYSQL_PS1="(\u@\h) [\d]> "
```

- *Use a command-line option.* You can set the --prompt option on the command line to mysql. For example:

```
shell> mysql --prompt="(\u@\h) [\d]> "
(user@host) [database]>
```

- *Use an option file.* You can set the prompt option in the [mysql] group of any MySQL option file, such as /etc/my.cnf or the .my.cnf file in your home directory. For example:

```
[mysql]
prompt=(\\u@\\h) [\\d]>\\_
```

In this example, note that the backslashes are doubled. If you set the prompt using the prompt option in an option file, it is advisable to double the backslashes when using the special prompt options. There is some overlap in the set of permissible prompt options and the set of special escape sequences that are recognized in option files. (The rules for escape sequences in option files are listed in Section 4.2.6, "Using Option Files".) The overlap may cause you problems if you use single backslashes. For example, \s is interpreted as a space rather than as the current seconds value. The following example shows how to define a prompt within an option file to include the current time in HH:MM:SS> format:

```
[mysql]
prompt="\\r:\\m:\\s> "
```

- *Set the prompt interactively.* You can change your prompt interactively by using the prompt (or \R) command. For example:

```
mysql> prompt (\u@\h) [\d]>\_
PROMPT set to '(\u@\h) [\d]>\_'
(user@host) [database]>
(user@host) [database]> prompt
Returning to default PROMPT of mysql>
mysql>
```

4.5.1.3 mysql Logging

The mysql client can do these types of logging for statements executed interactively:

- On Unix, mysql writes the statements to a history file. By default, this file is named .mysql_history in your home directory. To specify a different file, set the value of the MYSQL_HISTFILE environment variable.

- On all platforms, if the `--syslog` option is given, `mysql` writes the statements to the system logging facility. On Unix, this is `syslog`; on Windows, it is the Windows Event Log. The destination where logged messages appear is system dependent. On Linux, the destination is often the `/var/log/messages` file.

The following discussion describes characteristics that apply to all logging types and provides information specific to each logging type.

How Logging Occurs

For each enabled logging destination, statement logging occurs as follows:

- Statements are logged only when executed interactively. Statements are noninteractive, for example, when read from a file or a pipe. It is also possible to suppress statement logging by using the `--batch` or `--execute` option.
- Statements are ignored and not logged if they match any pattern in the “ignore” list. This list is described later.
- `mysql` logs each nonignored, nonempty statement line individually.
- If a nonignored statement spans multiple lines (not including the terminating delimiter), `mysql` concatenates the lines to form the complete statement, maps newlines to spaces, and logs the result, plus a delimiter.

Consequently, an input statement that spans multiple lines can be logged twice. Consider this input:

```
mysql> SELECT
    -> 'Today is'
    -> ,
    -> CURDATE()
    -> ;
```

In this case, `mysql` logs the “SELECT”, “'Today is'”, “,”, “CURDATE()”, and “;” lines as it reads them. It also logs the complete statement, after mapping `SELECT\n'Today is'\n,\nCURDATE()` to `SELECT 'Today is' , CURDATE()`, plus a delimiter. Thus, these lines appear in logged output:

```
SELECT
'Today is'
,
CURDATE()
;
SELECT 'Today is' , CURDATE();
```

`mysql` ignores for logging purposes statements that match any pattern in the “ignore” list. By default, the pattern list is `"*IDENTIFIED*:*PASSWORD*"`, to ignore statements that refer to passwords. Pattern matching is not case sensitive. Within patterns, two characters are special:

- `?` matches any single character.
- `*` matches any sequence of zero or more characters.

To specify additional patterns, use the `--histignore` option or set the `MYSQL_HISTIGNORE` environment variable. (If both are specified, the option value takes precedence.) The value should be a colon-separated list of one or more patterns, which are appended to the default pattern list.

Patterns specified on the command line might need to be quoted or escaped to prevent your command interpreter from treating them specially. For example, to suppress logging for `UPDATE` and `DELETE` statements in addition to statements that refer to passwords, invoke `mysql` like this:

```
shell> mysql --histignore="*UPDATE*:*DELETE*"
```

Controlling the History File

The .mysql_history file should be protected with a restrictive access mode because sensitive information might be written to it, such as the text of SQL statements that contain passwords. See Section 6.1.2.1, “End-User Guidelines for Password Security”.

If you do not want to maintain a history file, first remove .mysql_history if it exists. Then use either of the following techniques to prevent it from being created again:

- Set the MYSQL_HISTFILE environment variable to /dev/null. To cause this setting to take effect each time you log in, put it in one of your shell's startup files.
- Create .mysql_history as a symbolic link to /dev/null; this need be done only once:

```
shell> ln -s /dev/null $HOME/.mysql_history
```

syslog Logging Characteristics

If the --syslog option is given, mysql writes interactive statements to the system logging facility. Message logging has the following characteristics.

Logging occurs at the “information” level. This corresponds to the LOG_INFO priority for syslog on Unix/Linux syslog capability and to EVENTLOG_INFORMATION_TYPE for the Windows Event Log. Consult your system documentation for configuration of your logging capability.

Message size is limited to 1024 bytes.

Messages consist of the identifier MysqlClient followed by these values:

- SYSTEM_USER

  The system user name (login name) or -- if the user is unknown.

- MYSQL_USER

  The MySQL user name (specified with the --user option) or -- if the user is unknown.

- CONNECTION_ID:

  The client connection identifier. This is the same as the CONNECTION_ID() function value within the session.

- DB_SERVER

  The server host or -- if the host is unknown.

- DB

  The default database or -- if no database has been selected.

- QUERY

  The text of the logged statement.

Here is a sample of output generated on Linux by using --syslog. This output is formatted for readability; each logged message actually takes a single line.

```
Mar  7 12:39:25 myhost MysqlClient[20824]:
  SYSTEM_USER:'oscar', MYSQL_USER:'my_oscar', CONNECTION_ID:23,
  DB_SERVER:'127.0.0.1', DB:'--', QUERY:'USE test;'
Mar  7 12:39:28 myhost MysqlClient[20824]:
  SYSTEM_USER:'oscar', MYSQL_USER:'my_oscar', CONNECTION_ID:23,
  DB_SERVER:'127.0.0.1', DB:'test', QUERY:'SHOW TABLES;'
```

4.5.1.4 mysql Server-Side Help

```
mysql> help search_string
```

If you provide an argument to the `help` command, `mysql` uses it as a search string to access server-side help from the contents of the MySQL Reference Manual. The proper operation of this command requires that the help tables in the `mysql` database be initialized with help topic information (see Section 5.1.10, “Server-Side Help”).

If there is no match for the search string, the search fails:

```
mysql> help me

Nothing found
Please try to run 'help contents' for a list of all accessible topics
```

Use `help contents` to see a list of the help categories:

```
mysql> help contents
You asked for help about help category: "Contents"
For more information, type 'help <item>', where <item> is one of the
following categories:
   Account Management
   Administration
   Data Definition
   Data Manipulation
   Data Types
   Functions
   Functions and Modifiers for Use with GROUP BY
   Geographic Features
   Language Structure
   Plugins
   Storage Engines
   Stored Routines
   Table Maintenance
   Transactions
   Triggers
```

If the search string matches multiple items, `mysql` shows a list of matching topics:

```
mysql> help logs
Many help items for your request exist.
To make a more specific request, please type 'help <item>',
where <item> is one of the following topics:
   SHOW
   SHOW BINARY LOGS
   SHOW ENGINE
   SHOW LOGS
```

Use a topic as the search string to see the help entry for that topic:

```
mysql> help show binary logs
Name: 'SHOW BINARY LOGS'
Description:
Syntax:
SHOW BINARY LOGS
SHOW MASTER LOGS

Lists the binary log files on the server. This statement is used as
part of the procedure described in [purge-binary-logs], that shows how
to determine which logs can be purged.

mysql> SHOW BINARY LOGS;
+---------------+-----------+
| Log_name      | File_size |
+---------------+-----------+
```

```
| binlog.000015 |    724935 |
| binlog.000016 |    733481 |
+---------------+-----------+
```

The search string can contain the wildcard characters % and _. These have the same meaning as for pattern-matching operations performed with the LIKE operator. For example, HELP rep% returns a list of topics that begin with rep:

```
mysql> HELP rep%
Many help items for your request exist.
To make a more specific request, please type 'help <item>',
where <item> is one of the following
topics:
   REPAIR TABLE
   REPEAT FUNCTION
   REPEAT LOOP
   REPLACE
   REPLACE FUNCTION
```

4.5.1.5 Executing SQL Statements from a Text File

The mysql client typically is used interactively, like this:

```
shell> mysql db_name
```

However, it is also possible to put your SQL statements in a file and then tell mysql to read its input from that file. To do so, create a text file *text_file* that contains the statements you wish to execute. Then invoke mysql as shown here:

```
shell> mysql db_name < text_file
```

If you place a USE *db_name* statement as the first statement in the file, it is unnecessary to specify the database name on the command line:

```
shell> mysql < text_file
```

If you are already running mysql, you can execute an SQL script file using the source command or \. command:

```
mysql> source file_name
mysql> \. file_name
```

Sometimes you may want your script to display progress information to the user. For this you can insert statements like this:

```
SELECT '<info_to_display>' AS ' ';
```

The statement shown outputs <info_to_display>.

You can also invoke mysql with the --verbose option, which causes each statement to be displayed before the result that it produces.

mysql ignores Unicode byte order mark (BOM) characters at the beginning of input files. Previously, it read them and sent them to the server, resulting in a syntax error. Presence of a BOM does not cause mysql to change its default character set. To do that, invoke mysql with an option such as --default-character-set=utf8.

For more information about batch mode, see Section 3.5, “Using mysql in Batch Mode”.

4.5.1.6 mysql Tips

This section describes some techniques that can help you use mysql more effectively.

Input-Line Editing

mysql supports input-line editing, which enables you to modify the current input line in place or recall previous input lines. For example, the **left-arrow** and **right-arrow** keys move horizontally within the current input line, and the **up-arrow** and **down-arrow** keys move up and down through the set of previously entered lines. **Backspace** deletes the character before the cursor and typing new characters enters them at the cursor position. To enter the line, press **Enter**.

On Windows, the editing key sequences are the same as supported for command editing in console windows. On Unix, the key sequences depend on the input library used to build mysql (for example, the libedit or readline library).

Documentation for the libedit and readline libraries is available online. To change the set of key sequences permitted by a given input library, define key bindings in the library startup file. This is a file in your home directory: .editrc for libedit and .inputrc for readline.

For example, in libedit, **Control+W** deletes everything before the current cursor position and **Control+U** deletes the entire line. In readline, **Control+W** deletes the word before the cursor and **Control+U** deletes everything before the current cursor position. If mysql was built using libedit, a user who prefers the readline behavior for these two keys can put the following lines in the .editrc file (creating the file if necessary):

```
bind "^W" ed-delete-prev-word
bind "^U" vi-kill-line-prev
```

To see the current set of key bindings, temporarily put a line that says only bind at the end of .editrc. mysql will show the bindings when it starts.

Unicode Support on Windows

Windows provides APIs based on UTF-16LE for reading from and writing to the console; the mysql client for Windows is able to use these APIs. The Windows installer creates an item in the MySQL menu named MySQL command line client - Unicode. This item invokes the mysql client with properties set to communicate through the console to the MySQL server using Unicode.

To take advantage of this support manually, run mysql within a console that uses a compatible Unicode font and set the default character set to a Unicode character set that is supported for communication with the server:

1. Open a console window.
2. Go to the console window properties, select the font tab, and choose Lucida Console or some other compatible Unicode font. This is necessary because console windows start by default using a DOS raster font that is inadequate for Unicode.
3. Execute mysql.exe with the --default-character-set=utf8 (or utf8mb4) option. This option is necessary because utf16le is not supported as a connection character set.

With those changes, mysql will use the Windows APIs to communicate with the console using UTF-16LE, and communicate with the server using UTF-8. (The menu item mentioned previously sets the font and character set as just described.)

To avoid those steps each time you run mysql, you can create a shortcut that invokes mysql.exe. The shortcut should set the console font to Lucida Console or some other compatible Unicode font, and pass the --default-character-set=utf8 (or utf8mb4) option to mysql.exe.

Alternatively, create a shortcut that only sets the console font, and set the character set in the [mysql] group of your my.ini file:

```
[mysql]
default-character-set=utf8
```

Displaying Query Results Vertically

Some query results are much more readable when displayed vertically, instead of in the usual horizontal table format. Queries can be displayed vertically by terminating the query with \G instead of a semicolon. For example, longer text values that include newlines often are much easier to read with vertical output:

```
mysql> SELECT * FROM mails WHERE LENGTH(txt) < 300 LIMIT 300,1\G
*************************** 1. row ***************************
  msg_nro: 3068
     date: 2000-03-01 23:29:50
time_zone: +0200
mail_from: Monty
    reply: monty@no.spam.com
  mail_to: "Thimble Smith" <tim@no.spam.com>
      sbj: UTF-8
      txt: >>>>> "Thimble" == Thimble Smith writes:

Thimble> Hi.  I think this is a good idea.  Is anyone familiar
Thimble> with UTF-8 or Unicode? Otherwise, I'll put this on my
Thimble> TODO list and see what happens.

Yes, please do that.

Regards,
Monty
     file: inbox-jani-1
     hash: 190402944
1 row in set (0.09 sec)
```

Using the --safe-updates Option

For beginners, a useful startup option is `--safe-updates` (or `--i-am-a-dummy`, which has the same effect). It is helpful for cases when you might have issued a `DELETE FROM tbl_name` statement but forgotten the `WHERE` clause. Normally, such a statement deletes all rows from the table. With `--safe-updates`, you can delete rows only by specifying the key values that identify them. This helps prevent accidents.

When you use the `--safe-updates` option, `mysql` issues the following statement when it connects to the MySQL server:

```
SET sql_safe_updates=1, sql_select_limit=1000, max_join_size=1000000;
```

See Section 5.1.5, “Server System Variables”.

The `SET` statement has the following effects:

- You are not permitted to execute an `UPDATE` or `DELETE` statement unless you specify a key constraint in the `WHERE` clause or provide a `LIMIT` clause (or both). For example:

  ```
  UPDATE tbl_name SET not_key_column=val WHERE key_column=val;

  UPDATE tbl_name SET not_key_column=val LIMIT 1;
  ```

- The server limits all large `SELECT` results to 1,000 rows unless the statement includes a `LIMIT` clause.
- The server aborts multiple-table `SELECT` statements that probably need to examine more than 1,000,000 row combinations.

To specify limits different from 1,000 and 1,000,000, you can override the defaults by using the `--select_limit` and `--max_join_size` options:

```
shell> mysql --safe-updates --select_limit=500 --max_join_size=10000
```

Disabling mysql Auto-Reconnect

If the mysql client loses its connection to the server while sending a statement, it immediately and automatically tries to reconnect once to the server and send the statement again. However, even if mysql succeeds in reconnecting, your first connection has ended and all your previous session objects and settings are lost: temporary tables, the autocommit mode, and user-defined and session variables. Also, any current transaction rolls back. This behavior may be dangerous for you, as in the following example where the server was shut down and restarted between the first and second statements without you knowing it:

```
mysql> SET @a=1;
Query OK, 0 rows affected (0.05 sec)

mysql> INSERT INTO t VALUES(@a);
ERROR 2006: MySQL server has gone away
No connection. Trying to reconnect...
Connection id:    1
Current database: test

Query OK, 1 row affected (1.30 sec)

mysql> SELECT * FROM t;
+------+
| a    |
+------+
| NULL |
+------+
1 row in set (0.05 sec)
```

The @a user variable has been lost with the connection, and after the reconnection it is undefined. If it is important to have mysql terminate with an error if the connection has been lost, you can start the mysql client with the --skip-reconnect option.

For more information about auto-reconnect and its effect on state information when a reconnection occurs, see Section 27.7.19, “Controlling Automatic Reconnection Behavior”.

4.5.2 mysqladmin — Client for Administering a MySQL Server

mysqladmin is a client for performing administrative operations. You can use it to check the server's configuration and current status, to create and drop databases, and more.

Invoke mysqladmin like this:

```
shell> mysqladmin [options] command [command-arg] [command [command-arg]] ...
```

mysqladmin supports the following commands. Some of the commands take an argument following the command name.

- create *db_name*

  Create a new database named *db_name*.

- debug

  Tell the server to write debug information to the error log. Format and content of this information is subject to change.

  This includes information about the Event Scheduler. See Section 23.4.5, “Event Scheduler Status”.

- drop *db_name*

  Delete the database named *db_name* and all its tables.

- extended-status

Display the server status variables and their values.

- `flush-hosts`

  Flush all information in the host cache.

- `flush-logs [log_type ...]`

  Flush all logs.

  The `mysqladmin flush-logs` command permits optional log types to be given, to specify which logs to flush. Following the `flush-logs` command, you can provide a space-separated list of one or more of the following log types: `binary`, `engine`, `error`, `general`, `relay`, `slow`. These correspond to the log types that can be specified for the `FLUSH LOGS` SQL statement.

- `flush-privileges`

  Reload the grant tables (same as `reload`).

- `flush-status`

  Clear status variables.

- `flush-tables`

  Flush all tables.

- `flush-threads`

  Flush the thread cache.

- `kill id,id,...`

  Kill server threads. If multiple thread ID values are given, there must be no spaces in the list.

- `password new_password`

  Set a new password. This changes the password to *`new_password`* for the account that you use with `mysqladmin` for connecting to the server. Thus, the next time you invoke `mysqladmin` (or any other client program) using the same account, you will need to specify the new password.

  If the *`new_password`* value contains spaces or other characters that are special to your command interpreter, you need to enclose it within quotation marks. On Windows, be sure to use double quotation marks rather than single quotation marks; single quotation marks are not stripped from the password, but rather are interpreted as part of the password. For example:

```
shell> mysqladmin password "my new password"
```

  In MySQL 8.0, the new password can be omitted following the `password` command. In this case, `mysqladmin` prompts for the password value, which enables you to avoid specifying the password on the command line. Omitting the password value should be done only if `password` is the final command on the `mysqladmin` command line. Otherwise, the next argument is taken as the password.

Caution

Do not use this command used if the server was started with the `--skip-grant-tables` option. No password change will be applied. This is true even if you precede the `password` command with `flush-privileges` on the same command line to re-enable the grant tables because the flush operation occurs after you connect. However, you can use `mysqladmin`

> `flush-privileges` to re-enable the grant table and then use a separate `mysqladmin password` command to change the password.

- `ping`

  Check whether the server is available. The return status from `mysqladmin` is 0 if the server is running, 1 if it is not. This is 0 even in case of an error such as `Access denied`, because this means that the server is running but refused the connection, which is different from the server not running.

- `processlist`

  Show a list of active server threads. This is like the output of the `SHOW PROCESSLIST` statement. If the `--verbose` option is given, the output is like that of `SHOW FULL PROCESSLIST`. (See Section 13.7.5.29, “SHOW PROCESSLIST Syntax”.)

- `reload`

  Reload the grant tables.

- `refresh`

  Flush all tables and close and open log files.

- `shutdown`

  Stop the server.

- `start-slave`

  Start replication on a slave server.

- `status`

  Display a short server status message.

- `stop-slave`

  Stop replication on a slave server.

- `variables`

  Display the server system variables and their values.

- `version`

  Display version information from the server.

All commands can be shortened to any unique prefix. For example:

```
shell> mysqladmin proc stat
+----+-------+-----------+----+---------+------+-------+------------------+
| Id | User  | Host      | db | Command | Time | State | Info             |
+----+-------+-----------+----+---------+------+-------+------------------+
| 51 | monty | localhost |    | Query   | 0    |       | show processlist |
+----+-------+-----------+----+---------+------+-------+------------------+
Uptime: 1473624  Threads: 1  Questions: 39487
Slow queries: 0  Opens: 541  Flush tables: 1
Open tables: 19  Queries per second avg: 0.0268
```

The `mysqladmin status` command result displays the following values:

- `Uptime`

  The number of seconds the MySQL server has been running.

- `Threads`

  The number of active threads (clients).

- `Questions`

  The number of questions (queries) from clients since the server was started.

- `Slow queries`

  The number of queries that have taken more than `long_query_time` seconds. See Section 5.4.5, “The Slow Query Log”.

- `Opens`

  The number of tables the server has opened.

- `Flush tables`

  The number of `flush-*`, `refresh`, and `reload` commands the server has executed.

- `Open tables`

  The number of tables that currently are open.

If you execute `mysqladmin shutdown` when connecting to a local server using a Unix socket file, `mysqladmin` waits until the server's process ID file has been removed, to ensure that the server has stopped properly.

`mysqladmin` supports the following options, which can be specified on the command line or in the `[mysqladmin]` and `[client]` groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.9 `mysqladmin` Options

Format	Description
--bind-address	Use specified network interface to connect to MySQL Server
--compress	Compress all information sent between client and server
--connect_timeout	Number of seconds before connection timeout
--count	Number of iterations to make for repeated command execution
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--default-character-set	Specify default character set
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--enable-cleartext-plugin	Enable cleartext authentication plugin
--force	Continue even if an SQL error occurs
--help	Display help message and exit
--host	Connect to MySQL server on given host
--login-path	Read login path options from .mylogin.cnf
--no-beep	Do not beep when errors occur

Format	Description
--no-defaults	Read no option files
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use
--relative	Show the difference between the current and previous values when used with the --sleep option
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--show-warnings	Show warnings after statement execution
--shutdown_timeout	The maximum number of seconds to wait for server shutdown
--silent	Silent mode
--sleep	Execute commands repeatedly, sleeping for delay seconds in between
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tls-version	Protocols permitted for secure connections
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version	Display version information and exit
--vertical	Print query output rows vertically (one line per column value)
--wait	If the connection cannot be established, wait and retry instead of aborting

- `--help`, `-?`

  Display a help message and exit.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--count=N`, `-c N`

  The number of iterations to make for repeated command execution if the `--sleep` option is given.

- `--debug[=debug_options]`, `-# [debug_options]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:o,/tmp/mysqladmin.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, “Character Set Configuration”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqladmin` normally reads the `[client]` and `[mysqladmin]` groups. If the `--defaults-group-suffix=_other` option is given, `mysqladmin` also reads the `[client_other]` and `[mysqladmin_other]` groups.

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--force`, `-f`

  Do not ask for confirmation for the `drop db_name` command. With multiple commands, continue even if an error occurs.

- `--host=host_name`, `-h host_name`

  Connect to the MySQL server on the given host.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

- `--no-beep`, `-b`

  Suppress the warning beep that is emitted by default for errors such as a failure to connect to the server.

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.)

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *password* value following the `--password` or `-p` option on the command line, `mysqladmin` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqladmin` does not find it. See Section 6.3.10, “Pluggable Authentication”.

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--relative`, `-r`

  Show the difference between the current and previous values when used with the `--sleep` option. This option works only with the `extended-status` command.

- `--show-warnings`

Show warnings resulting from execution of statements sent to the server.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--shared-memory-base-name=`*`name`*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--silent`, `-s`

  Exit silently if a connection to the server cannot be established.

- `--sleep=`*`delay`*, `-i` *`delay`*

  Execute commands repeatedly, sleeping for *`delay`* seconds in between. The `--count` option determines the number of iterations. If `--count` is not given, `mysqladmin` executes commands indefinitely until interrupted.

- `--socket=`*`path`*, `-S` *`path`*

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=`*`protocol_list`*

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=`*`user_name`*, `-u` *`user_name`*

  The MySQL user name to use when connecting to the server.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does.

- `--version`, `-V`

Display version information and exit.

- `--vertical`, `-E`

  Print output vertically. This is similar to `--relative`, but prints output vertically.

- `--wait[=count]`, `-w[count]`

  If the connection cannot be established, wait and retry instead of aborting. If a *`count`* value is given, it indicates the number of times to retry. The default is one time.

You can also set the following variables by using `--var_name=value`.

- `connect_timeout`

  The maximum number of seconds before connection timeout. The default value is 43200 (12 hours).

- `shutdown_timeout`

  The maximum number of seconds to wait for server shutdown. The default value is 3600 (1 hour).

4.5.3 mysqlcheck — A Table Maintenance Program

The `mysqlcheck` client performs table maintenance: It checks, repairs, optimizes, or analyzes tables.

Each table is locked and therefore unavailable to other sessions while it is being processed, although for check operations, the table is locked with a `READ` lock only (see Section 13.3.5, “LOCK TABLES and UNLOCK TABLES Syntax”, for more information about `READ` and `WRITE` locks). Table maintenance operations can be time-consuming, particularly for large tables. If you use the `--databases` or `--all-databases` option to process all tables in one or more databases, an invocation of `mysqlcheck` might take a long time. (This is also true for `mysql_upgrade` because that program invokes `mysqlcheck` to check all tables and repair them if necessary.)

`mysqlcheck` is similar in function to `myisamchk`, but works differently. The main operational difference is that `mysqlcheck` must be used when the `mysqld` server is running, whereas `myisamchk` should be used when it is not. The benefit of using `mysqlcheck` is that you do not have to stop the server to perform table maintenance.

`mysqlcheck` uses the SQL statements `CHECK TABLE`, `REPAIR TABLE`, `ANALYZE TABLE`, and `OPTIMIZE TABLE` in a convenient way for the user. It determines which statements to use for the operation you want to perform, and then sends the statements to the server to be executed. For details about which storage engines each statement works with, see the descriptions for those statements in Section 13.7.2, “Table Maintenance Statements”.

The `MyISAM` storage engine supports all four maintenance operations, so `mysqlcheck` can be used to perform any of them on `MyISAM` tables. Other storage engines do not necessarily support all operations. In such cases, an error message is displayed. For example, if `test.t` is a `MEMORY` table, an attempt to check it produces this result:

```
shell> mysqlcheck test t
test.t
note     : The storage engine for the table doesn't support check
```

If `mysqlcheck` is unable to repair a table, see Section 2.10.3, “Rebuilding or Repairing Tables or Indexes” for manual table repair strategies. This will be the case, for example, for `InnoDB` tables, which can be checked with `CHECK TABLE`, but not repaired with `REPAIR TABLE`.

Caution

It is best to make a backup of a table before performing a table repair operation; under some circumstances the operation might cause data loss. Possible causes include but are not limited to file system errors.

There are three general ways to invoke mysqlcheck:

```
shell> mysqlcheck [options] db_name [tbl_name ...]
shell> mysqlcheck [options] --databases db_name ...
shell> mysqlcheck [options] --all-databases
```

If you do not name any tables following *db_name* or if you use the --databases or --all-databases option, entire databases are checked.

mysqlcheck has a special feature compared to other client programs. The default behavior of checking tables (--check) can be changed by renaming the binary. If you want to have a tool that repairs tables by default, you should just make a copy of mysqlcheck named mysqlrepair, or make a symbolic link to mysqlcheck named mysqlrepair. If you invoke mysqlrepair, it repairs tables.

The names shown in the following table can be used to change mysqlcheck default behavior.

Command	Meaning
mysqlrepair	The default option is --repair
mysqlanalyze	The default option is --analyze
mysqloptimize	The default option is --optimize

mysqlcheck supports the following options, which can be specified on the command line or in the [mysqlcheck] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.10 mysqlcheck Options

Format	Description
--all-databases	Check all tables in all databases
--all-in-1	Execute a single statement for each database that names all the tables from that database
--analyze	Analyze the tables
--auto-repair	If a checked table is corrupted, automatically fix it
--bind-address	Use specified network interface to connect to MySQL Server
--character-sets-dir	Directory where character sets are installed
--check	Check the tables for errors
--check-only-changed	Check only tables that have changed since the last check
--check-upgrade	Invoke CHECK TABLE with the FOR UPGRADE option
--compress	Compress all information sent between client and server
--databases	Interpret all arguments as database names
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--default-character-set	Specify default character set
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--enable-cleartext-plugin	Enable cleartext authentication plugin
--extended	Check and repair tables

Format	Description
--fast	Check only tables that have not been closed properly
--force	Continue even if an SQL error occurs
--help	Display help message and exit
--host	Connect to MySQL server on given host
--login-path	Read login path options from .mylogin.cnf
--medium-check	Do a check that is faster than an --extended operation
--no-defaults	Read no option files
--optimize	Optimize the tables
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use
--quick	The fastest method of checking
--repair	Perform a repair that can fix almost anything except unique keys that are not unique
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--silent	Silent mode
--skip-database	Omit this database from performed operations
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tables	Overrides the --databases or -B option
--tls-version	Protocols permitted for secure connections
--use-frm	For repair operations on MyISAM tables
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version	Display version information and exit
--write-binlog	Log ANALYZE, OPTIMIZE, REPAIR statements to binary log. --skip-write-binlog adds NO_WRITE_TO_BINLOG to these statements.

- `--help`, `-?`

Display a help message and exit.

- `--all-databases`, `-A`

  Check all tables in all databases. This is the same as using the `--databases` option and naming all the databases on the command line, except that the `INFORMATION_SCHEMA` and `performace_schema` databases are not checked. They can be checked by explicitly naming them with the `--databases` option.

- `--all-in-1`, `-1`

  Instead of issuing a statement for each table, execute a single statement for each database that names all the tables from that database to be processed.

- `--analyze`, `-a`

  Analyze the tables.

- `--auto-repair`

  If a checked table is corrupted, automatically fix it. Any necessary repairs are done after all tables have been checked.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--check`, `-c`

  Check the tables for errors. This is the default operation.

- `--check-only-changed`, `-C`

  Check only tables that have changed since the last check or that have not been closed properly.

- `--check-upgrade`, `-g`

  Invoke `CHECK TABLE` with the `FOR UPGRADE` option to check tables for incompatibilities with the current version of the server.

- `--compress`

  Compress all information sent between the client and the server if both support compression.

- `--databases`, `-B`

  Process all tables in the named databases. Normally, `mysqlcheck` treats the first name argument on the command line as a database name and any following names as table names. With this option, it treats all name arguments as database names.

- `--debug[=`*`debug_options`*`]`, `-# [`*`debug_options`*`]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:o`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, “Character Set Configuration”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqlcheck` normally reads the `[client]` and `[mysqlcheck]` groups. If the `--defaults-group-suffix=_other` option is given, `mysqlcheck` also reads the `[client_other]` and `[mysqlcheck_other]` groups.

- `--extended`, `-e`

  If you are using this option to check tables, it ensures that they are 100% consistent but takes a long time.

  If you are using this option to repair tables, it runs an extended repair that may not only take a long time to execute, but may produce a lot of garbage rows also!

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--fast`, `-F`

  Check only tables that have not been closed properly.

- `--force`, `-f`

  Continue even if an SQL error occurs.

- `--host=host_name`, `-h host_name`

  Connect to the MySQL server on the given host.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to

authenticate as. To create or modify a login path file, use the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

- --medium-check, -m

  Do a check that is faster than an --extended operation. This finds only 99.99% of all errors, which should be good enough in most cases.

- --no-defaults

  Do not read any option files. If program startup fails due to reading unknown options from an option file, --no-defaults can be used to prevent them from being read.

  The exception is that the .mylogin.cnf file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when --no-defaults is used. (.mylogin.cnf is created by the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.)

- --optimize, -o

  Optimize the tables.

- --password[=*password*], -p[*password*]

  The password to use when connecting to the server. If you use the short option form (-p), you *cannot* have a space between the option and the password. If you omit the *password* value following the --password or -p option on the command line, mysqlcheck prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- --pipe, -W

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- --plugin-dir=*dir_name*

  The directory in which to look for plugins. Specify this option if the --default-auth option is used to specify an authentication plugin but mysqlcheck does not find it. See Section 6.3.10, “Pluggable Authentication”.

- --port=*port_num*, -P *port_num*

  The TCP/IP port number to use for the connection.

- --print-defaults

  Print the program name and all options that it gets from option files.

- --protocol={TCP|SOCKET|PIPE|MEMORY}

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- --quick, -q

  If you are using this option to check tables, it prevents the check from scanning the rows to check for incorrect links. This is the fastest check method.

If you are using this option to repair tables, it tries to repair only the index tree. This is the fastest repair method.

- `--repair`, `-r`

  Perform a repair that can fix almost anything except unique keys that are not unique.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--shared-memory-base-name=`*`name`*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--silent`, `-s`

  Silent mode. Print only error messages.

- `--skip-database=`*`db_name`*

  Do not include the named database (case sensitive) in the operations performed by `mysqlcheck`.

- `--socket=`*`path`*, `-S` *`path`*

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tables`

  Override the `--databases` or `-B` option. All name arguments following the option are regarded as table names.

- `--tls-version=`*`protocol_list`*

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--use-frm`

For repair operations on MyISAM tables, get the table structure from the data dictionary so that the table can be repaired even if the .MYI header is corrupted.

- --user=*user_name*, -u *user_name*

  The MySQL user name to use when connecting to the server.

- --verbose, -v

  Verbose mode. Print information about the various stages of program operation.

- --version, -V

  Display version information and exit.

- --write-binlog

  This option is enabled by default, so that ANALYZE TABLE, OPTIMIZE TABLE, and REPAIR TABLE statements generated by mysqlcheck are written to the binary log. Use --skip-write-binlog to cause NO_WRITE_TO_BINLOG to be added to the statements so that they are not logged. Use the --skip-write-binlog when these statements should not be sent to replication slaves or run when using the binary logs for recovery from backup.

4.5.4 mysqldump — A Database Backup Program

The mysqldump client utility performs logical backups, producing a set of SQL statements that can be executed to reproduce the original database object definitions and table data. It dumps one or more MySQL databases for backup or transfer to another SQL server. The mysqldump command can also generate output in CSV, other delimited text, or XML format.

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mysqldump requires at least the SELECT privilege for dumped tables, SHOW VIEW for dumped views, TRIGGER for dumped triggers, and LOCK TABLES if the --single-transaction option is not used. Certain options might require other privileges as noted in the option descriptions.

To reload a dump file, you must have the privileges required to execute the statements that it contains, such as the appropriate CREATE privileges for objects created by those statements.

mysqldump output can include ALTER DATABASE statements that change the database collation. These may be used when dumping stored programs to preserve their character encodings. To reload a dump file containing such statements, the ALTER privilege for the affected database is required.

Note

A dump made using PowerShell on Windows with output redirection creates a file that has UTF-16 encoding:

```
shell> mysqldump [options] > dump.sql
```

However, UTF-16 is not permitted as a connection character set (see Section 10.1.4, “Connection Character Sets and Collations”), so the dump file will not load correctly. To work around this issue, use the --result-file option, which creates the output in ASCII format:

```
shell> mysqldump [options] --result-file=dump.sql
```

Performance and Scalability Considerations

mysqldump advantages include the convenience and flexibility of viewing or even editing the output before restoring. You can clone databases for development and DBA work, or produce slight variations of an existing database for testing. It is not intended as a fast or scalable solution for backing up substantial amounts of data. With large data sizes, even if the backup step takes a reasonable time, restoring the data can be very slow because replaying the SQL statements involves disk I/O for insertion, index creation, and so on.

For large-scale backup and restore, a physical backup is more appropriate, to copy the data files in their original format that can be restored quickly:

- If your tables are primarily InnoDB tables, or if you have a mix of InnoDB and MyISAM tables, consider using the mysqlbackup command of the MySQL Enterprise Backup product. (Available as part of the Enterprise subscription.) It provides the best performance for InnoDB backups with minimal disruption; it can also back up tables from MyISAM and other storage engines; and it provides a number of convenient options to accommodate different backup scenarios. See Section 29.2, “MySQL Enterprise Backup Overview”.

mysqldump can retrieve and dump table contents row by row, or it can retrieve the entire content from a table and buffer it in memory before dumping it. Buffering in memory can be a problem if you are dumping large tables. To dump tables row by row, use the --quick option (or --opt, which enables --quick). The --opt option (and hence --quick) is enabled by default, so to enable memory buffering, use --skip-quick.

If you are using a recent version of mysqldump to generate a dump to be reloaded into a very old MySQL server, use the --skip-opt option instead of the --opt or --extended-insert option.

For additional information about mysqldump, see Section 7.4, “Using mysqldump for Backups”.

Invocation Syntax

There are in general three ways to use mysqldump—in order to dump a set of one or more tables, a set of one or more complete databases, or an entire MySQL server—as shown here:

```
shell> mysqldump [options] db_name [tbl_name ...]
shell> mysqldump [options] --databases db_name ...
shell> mysqldump [options] --all-databases
```

To dump entire databases, do not name any tables following *db_name*, or use the `--databases` or `--all-databases` option.

To see a list of the options your version of `mysqldump` supports, issue the command `mysqldump --help`.

Option Syntax - Alphabetical Summary

`mysqldump` supports the following options, which can be specified on the command line or in the `[mysqldump]` and `[client]` groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.11 `mysqldump` Options

Format	Description	Introduced
--add-drop-database	Add DROP DATABASE statement before each CREATE DATABASE statement	
--add-drop-table	Add DROP TABLE statement before each CREATE TABLE statement	
--add-drop-trigger	Add DROP TRIGGER statement before each CREATE TRIGGER statement	
--add-locks	Surround each table dump with LOCK TABLES and UNLOCK TABLES statements	
--all-databases	Dump all tables in all databases	
--allow-keywords	Allow creation of column names that are keywords	
--apply-slave-statements	Include STOP SLAVE prior to CHANGE MASTER statement and START SLAVE at end of output	
--bind-address	Use specified network interface to connect to MySQL Server	
--character-sets-dir	Directory where character sets are installed	
--column-statistics	Write ANALYZE TABLE statements to generate statistics histograms	8.0.2
--comments	Add comments to dump file	
--compact	Produce more compact output	
--compatible	Produce output that is more compatible with other database systems or with older MySQL servers	
--complete-insert	Use complete INSERT statements that include column names	
--compress	Compress all information sent between client and server	
--create-options	Include all MySQL-specific table options in CREATE TABLE statements	
--databases	Interpret all name arguments as database names	
--debug	Write debugging log	
--debug-check	Print debugging information when program exits	
--debug-info	Print debugging information, memory, and CPU statistics when program exits	
--default-auth	Authentication plugin to use	
--default-character-set	Specify default character set	

Format	Description	Introduced
--defaults-extra-file	Read named option file in addition to usual option files	
--defaults-file	Read only named option file	
--defaults-group-suffix	Option group suffix value	
--delete-master-logs	On a master replication server, delete the binary logs after performing the dump operation	
--disable-keys	For each table, surround INSERT statements with statements to disable and enable keys	
--dump-date	Include dump date as "Dump completed on" comment if --comments is given	
--dump-slave	Include CHANGE MASTER statement that lists binary log coordinates of slave's master	
--enable-cleartext-plugin	Enable cleartext authentication plugin	
--events	Dump events from dumped databases	
--extended-insert	Use multiple-row INSERT syntax	
--fields-enclosed-by	This option is used with the --tab option and has the same meaning as the corresponding clause for LOAD DATA INFILE	
--fields-escaped-by	This option is used with the --tab option and has the same meaning as the corresponding clause for LOAD DATA INFILE	
--fields-optionally-enclosed-by	This option is used with the --tab option and has the same meaning as the corresponding clause for LOAD DATA INFILE	
--fields-terminated-by	This option is used with the --tab option and has the same meaning as the corresponding clause for LOAD DATA INFILE	
--flush-logs	Flush MySQL server log files before starting dump	
--flush-privileges	Emit a FLUSH PRIVILEGES statement after dumping mysql database	
--force	Continue even if an SQL error occurs during a table dump	
--help	Display help message and exit	
--hex-blob	Dump binary columns using hexadecimal notation	
--host	Host to connect to (IP address or hostname)	
--ignore-error	Ignore specified errors	
--ignore-table	Do not dump given table	
--include-master-host-port	Include MASTER_HOST/MASTER_PORT options in CHANGE MASTER statement produced with --dump-slave	
--insert-ignore	Write INSERT IGNORE rather than INSERT statements	
--lines-terminated-by	This option is used with the --tab option and has the same meaning as the corresponding clause for LOAD DATA INFILE	
--lock-all-tables	Lock all tables across all databases	
--lock-tables	Lock all tables before dumping them	
--log-error	Append warnings and errors to named file	

Format	Description	Introduced
--login-path	Read login path options from .mylogin.cnf	
--master-data	Write the binary log file name and position to the output	
--max_allowed_packet	Maximum packet length to send to or receive from server	
--net_buffer_length	Buffer size for TCP/IP and socket communication	
--network-timeout	Increase network timeouts to permit larger table dumps	8.0.1
--no-autocommit	Enclose the INSERT statements for each dumped table within SET autocommit = 0 and COMMIT statements	
--no-create-db	Do not write CREATE DATABASE statements	
--no-create-info	Do not write CREATE TABLE statements that re-create each dumped table	
--no-data	Do not dump table contents	
--no-defaults	Read no option files	
--no-set-names	Same as --skip-set-charset	
--no-tablespaces	Do not write any CREATE LOGFILE GROUP or CREATE TABLESPACE statements in output	
--opt	Shorthand for --add-drop-table --add-locks --create-options --disable-keys --extended-insert --lock-tables --quick --set-charset.	
--order-by-primary	Dump each table's rows sorted by its primary key, or by its first unique index	
--password	Password to use when connecting to server	
--pipe	On Windows, connect to server using named pipe	
--plugin-dir	Directory where plugins are installed	
--port	TCP/IP port number to use for connection	
--print-defaults	Print default options	
--protocol	Connection protocol to use	
--quick	Retrieve rows for a table from the server a row at a time	
--quote-names	Quote identifiers within backtick characters	
--replace	Write REPLACE statements rather than INSERT statements	
--result-file	Direct output to a given file	
--routines	Dump stored routines (procedures and functions) from dumped databases	
--secure-auth	Do not send passwords to server in old (pre-4.1) format	
--set-charset	Add SET NAMES default_character_set to output	
--set-gtid-purged	Whether to add SET @@GLOBAL.GTID_PURGED to output	
--shared-memory-base-name	The name of shared memory to use for shared-memory connections	
--single-transaction	Issue a BEGIN SQL statement before dumping data from server	
--skip-add-drop-table	Do not add a DROP TABLE statement before each CREATE TABLE statement	
--skip-add-locks	Do not add locks	

Format	Description	Introduced
--skip-comments	Do not add comments to dump file	
--skip-compact	Do not produce more compact output	
--skip-disable-keys	Do not disable keys	
--skip-extended-insert	Turn off extended-insert	
--skip-opt	Turn off options set by --opt	
--skip-quick	Do not retrieve rows for a table from the server a row at a time	
--skip-quote-names	Do not quote identifiers	
--skip-set-charset	Do not write SET NAMES statement	
--skip-triggers	Do not dump triggers	
--skip-tz-utc	Turn off tz-utc	
--socket	For connections to localhost, the Unix socket file to use	
--ssl-ca	Path of file that contains list of trusted SSL CAs	
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format	
--ssl-cert	Path of file that contains X509 certificate in PEM format	
--ssl-cipher	List of permitted ciphers to use for connection encryption	
--ssl-crl	Path of file that contains certificate revocation lists	
--ssl-crlpath	Path of directory that contains certificate revocation list files	
--ssl-key	Path of file that contains X509 key in PEM format	
--ssl-mode	Security state of connection to server	
--tab	Produce tab-separated data files	
--tables	Override --databases or -B option	
--tls-version	Protocols permitted for secure connections	
--triggers	Dump triggers for each dumped table	
--tz-utc	Add SET TIME_ZONE='+00:00' to dump file	
--user	MySQL user name to use when connecting to server	
--verbose	Verbose mode	
--version	Display version information and exit	
--where	Dump only rows selected by given WHERE condition	
--xml	Produce XML output	

Connection Options

The `mysqldump` command logs into a MySQL server to extract information. The following options specify how to connect to the MySQL server, either on the same machine or a remote system.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--host=host_name`, `-h host_name`

  Dump data from the MySQL server on the given host. The default host is `localhost`.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--network-timeout`, `-M`

  Enable large tables to be dumped by setting `max_allowed_packet` to its maximum value and network read and write timeouts to a large value. This option is enabled by default. To disable it, use `--skip-network-timeout`.

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysqldump` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqldump` does not find it. See Section 6.3.10, “Pluggable Authentication”.

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

Note

Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=protocol_list`

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server.

You can also set the following variables by using `--var_name=value` syntax:

- `max_allowed_packet`

  The maximum size of the buffer for client/server communication. The default is 24MB, the maximum is 1GB.

- `net_buffer_length`

  The initial size of the buffer for client/server communication. When creating multiple-row `INSERT` statements (as with the `--extended-insert` or `--opt` option), `mysqldump` creates rows up to `net_buffer_length` bytes long. If you increase this variable, ensure that the MySQL server `net_buffer_length` system variable has a value at least this large.

Option-File Options

These options are used to control which option files to read.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. `file_name` is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. `file_name` is interpreted relative to the current directory if given as a relative path name rather than a full path name.

Exception: Even with --defaults-file, client programs read .mylogin.cnf.

- --defaults-group-suffix=*str*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *str*. For example, mysqldump normally reads the [client] and [mysqldump] groups. If the --defaults-group-suffix=_other option is given, mysqldump also reads the [client_other] and [mysqldump_other] groups.

- --no-defaults

  Do not read any option files. If program startup fails due to reading unknown options from an option file, --no-defaults can be used to prevent them from being read.

  The exception is that the .mylogin.cnf file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when --no-defaults is used. (.mylogin.cnf is created by the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.)

- --print-defaults

  Print the program name and all options that it gets from option files.

DDL Options

Usage scenarios for mysqldump include setting up an entire new MySQL instance (including database tables), and replacing data inside an existing instance with existing databases and tables. The following options let you specify which things to tear down and set up when restoring a dump, by encoding various DDL statements within the dump file.

- --add-drop-database

  Write a DROP DATABASE statement before each CREATE DATABASE statement. This option is typically used in conjunction with the --all-databases or --databases option because no CREATE DATABASE statements are written unless one of those options is specified.

- --add-drop-table

  Write a DROP TABLE statement before each CREATE TABLE statement.

- --add-drop-trigger

  Write a DROP TRIGGER statement before each CREATE TRIGGER statement.

- --all-tablespaces, -Y

  Adds to a table dump all SQL statements needed to create any tablespaces used by an NDB table. This information is not otherwise included in the output from mysqldump. This option is currently relevant only to MySQL Cluster tables, which are not supported in MySQL 8.0.

- --no-create-db, -n

  Suppress the CREATE DATABASE statements that are otherwise included in the output if the --databases or --all-databases option is given.

- --no-create-info, -t

  Do not write CREATE TABLE statements that create each dumped table.

Note

This option does *not* exclude statements creating log file groups or tablespaces from `mysqldump` output; however, you can use the `--no-tablespaces` option for this purpose.

- `--no-tablespaces`, `-y`

  This option suppresses all `CREATE LOGFILE GROUP` and `CREATE TABLESPACE` statements in the output of `mysqldump`.

- `--replace`

  Write `REPLACE` statements rather than `INSERT` statements.

Debug Options

The following options print debugging information, encode debugging information in the dump file, or let the dump operation proceed regardless of potential problems.

- `--allow-keywords`

  Permit creation of column names that are keywords. This works by prefixing each column name with the table name.

- `--comments`, `-i`

  Write additional information in the dump file such as program version, server version, and host. This option is enabled by default. To suppress this additional information, use `--skip-comments`.

- `--debug[=`*`debug_options`*`]`, `-# [`*`debug_options`*`]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default value is `d:t:o,/tmp/mysqldump.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--dump-date`

  If the `--comments` option is given, `mysqldump` produces a comment at the end of the dump of the following form:

  ```
  -- Dump completed on DATE
  ```

  However, the date causes dump files taken at different times to appear to be different, even if the data are otherwise identical. `--dump-date` and `--skip-dump-date` control whether the date is added to the comment. The default is `--dump-date` (include the date in the comment). `--skip-dump-date` suppresses date printing.

- `--force`, `-f`

  Ignore all errors; continue even if an SQL error occurs during a table dump.

  One use for this option is to cause `mysqldump` to continue executing even when it encounters a view that has become invalid because the definition refers to a table that has been dropped. Without `--force`, `mysqldump` exits with an error message. With `--force`, `mysqldump` prints the error

message, but it also writes an SQL comment containing the view definition to the dump output and continues executing.

If the `--ignore-error` option is also given to ignore specific errors, `--force` takes precedence.

- `--log-error=file_name`

  Log warnings and errors by appending them to the named file. The default is to do no logging.

- `--skip-comments`

  See the description for the `--comments` option.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does.

Help Options

The following options display information about the `mysqldump` command itself.

- `--help`, `-?`

  Display a help message and exit.

- `--version`, `-V`

  Display version information and exit.

Internationalization Options

The following options change how the `mysqldump` command represents character data with national language settings.

- `--character-sets-dir=dir_name`

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--default-character-set=charset_name`

  Use `charset_name` as the default character set. See Section 10.5, “Character Set Configuration”. If no character set is specified, `mysqldump` uses `utf8`.

- `--no-set-names`, `-N`

  Turns off the `--set-charset` setting, the same as specifying `--skip-set-charset`.

- `--set-charset`

  Write `SET NAMES default_character_set` to the output. This option is enabled by default. To suppress the `SET NAMES` statement, use `--skip-set-charset`.

Replication Options

The `mysqldump` command is frequently used to create an empty instance, or an instance including data, on a slave server in a replication configuration. The following options apply to dumping and restoring data on replication master and slave servers.

- `--apply-slave-statements`

  For a slave dump produced with the `--dump-slave` option, add a `STOP SLAVE` statement before the `CHANGE MASTER TO` statement and a `START SLAVE` statement at the end of the output.

- `--delete-master-logs`

On a master replication server, delete the binary logs by sending a PURGE BINARY LOGS statement to the server after performing the dump operation. This option automatically enables --master-data.

- --dump-slave[=*value*]

  This option is similar to --master-data except that it is used to dump a replication slave server to produce a dump file that can be used to set up another server as a slave that has the same master as the dumped server. It causes the dump output to include a CHANGE MASTER TO statement that indicates the binary log coordinates (file name and position) of the dumped slave's master. The CHANGE MASTER TO statement reads the values of Relay_Master_Log_File and Exec_Master_Log_Pos from the SHOW SLAVE STATUS output and uses them for MASTER_LOG_FILE and MASTER_LOG_POS respectively. These are the master server coordinates from which the slave should start replicating.

  Note

  Inconsistencies in the sequence of transactions from the relay log which have been executed can cause the wrong position to be used. See Section 18.4.1.34, "Replication and Transaction Inconsistencies" for more information.

  --dump-slave causes the coordinates from the master to be used rather than those of the dumped server, as is done by the --master-data option. In addition, specfiying this option causes the --master-data option to be overridden, if used, and effectively ignored.

  Warning

  This option should not be used if the server where the dump is going to be applied uses gtid_mode=ON and MASTER_AUTOPOSITION=1.

  The option value is handled the same way as for --master-data (setting no value or 1 causes a CHANGE MASTER TO statement to be written to the dump, setting 2 causes the statement to be written but encased in SQL comments) and has the same effect as --master-data in terms of enabling or disabling other options and in how locking is handled.

  This option causes mysqldump to stop the slave SQL thread before the dump and restart it again after.

  In conjunction with --dump-slave, the --apply-slave-statements and --include-master-host-port options can also be used.

- --include-master-host-port

  For the CHANGE MASTER TO statement in a slave dump produced with the --dump-slave option, add MASTER_HOST and MASTER_PORT options for the host name and TCP/IP port number of the slave's master.

- --master-data[=*value*]

  Use this option to dump a master replication server to produce a dump file that can be used to set up another server as a slave of the master. It causes the dump output to include a CHANGE MASTER TO statement that indicates the binary log coordinates (file name and position) of the dumped server. These are the master server coordinates from which the slave should start replicating after you load the dump file into the slave.

  If the option value is 2, the CHANGE MASTER TO statement is written as an SQL comment, and thus is informative only; it has no effect when the dump file is reloaded. If the option value is 1, the statement is not written as a comment and takes effect when the dump file is reloaded. If no option value is specified, the default value is 1.

This option requires the RELOAD privilege and the binary log must be enabled.

The --master-data option automatically turns off --lock-tables. It also turns on --lock-all-tables, unless --single-transaction also is specified, in which case, a global read lock is acquired only for a short time at the beginning of the dump (see the description for --single-transaction). In all cases, any action on logs happens at the exact moment of the dump.

It is also possible to set up a slave by dumping an existing slave of the master, using the --dump-slave option, which overrides --master-data and causes it to be ignored if both options are used.

- --set-gtid-purged=*value*

  This option enables control over global transaction ID (GTID) information written to the dump file, by indicating whether to add a SET @@global.gtid_purged statement to the output. This option may also cause a statement to be written to the output that disables binary logging while the dump file is being reloaded.

  The following table shows the permitted option values. The default value is AUTO.

Value	Meaning
OFF	Add no SET statement to the output.
ON	Add a SET statement to the output. An error occurs if GTIDs are not enabled on the server.
AUTO	Add a SET statement to the output if GTIDs are enabled on the server.

  The --set-gtid-purged option has the following effect on binary logging when the dump file is reloaded:

  - --set-gtid-purged=OFF: SET @@SESSION.SQL_LOG_BIN=0; is not added to the output.
  - --set-gtid-purged=ON: SET @@SESSION.SQL_LOG_BIN=0; is added to the output.
  - --set-gtid-purged=AUTO: SET @@SESSION.SQL_LOG_BIN=0; is added to the output if GTIDs are enabled on the server you are backing up (that is, if AUTO evaluates to ON).

Format Options

The following options specify how to represent the entire dump file or certain kinds of data in the dump file. They also control whether certain optional information is written to the dump file.

- --compact

  Produce more compact output. This option enables the --skip-add-drop-table, --skip-add-locks, --skip-comments, --skip-disable-keys, and --skip-set-charset options.

- --compatible=*name*

  Produce output that is more compatible with other database systems or with older MySQL servers. The value of *name* can be ansi, mysql323, mysql40, postgresql, oracle, mssql, db2, maxdb, no_key_options, no_table_options, or no_field_options. To use several values, separate them by commas. These values have the same meaning as the corresponding options for setting the server SQL mode. See Section 5.1.8, “Server SQL Modes”.

  This option does not guarantee compatibility with other servers. It only enables those SQL mode values that are currently available for making dump output more compatible. For example, --compatible=oracle does not map data types to Oracle types or use Oracle comment syntax.

  This option requires a server version of 4.1.0 or higher. With older servers, it does nothing.

- `--complete-insert`, `-c`

  Use complete `INSERT` statements that include column names.

- `--create-options`

  Include all MySQL-specific table options in the `CREATE TABLE` statements.

- `--fields-terminated-by=...`, `--fields-enclosed-by=...`, `--fields-optionally-enclosed-by=...`, `--fields-escaped-by=...`

  These options are used with the `--tab` option and have the same meaning as the corresponding `FIELDS` clauses for `LOAD DATA INFILE`. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

- `--hex-blob`

  Dump binary columns using hexadecimal notation (for example, `'abc'` becomes `0x616263`). The affected data types are `BINARY`, `VARBINARY`, the `BLOB` types, and `BIT`.

- `--lines-terminated-by=...`

  This option is used with the `--tab` option and has the same meaning as the corresponding `LINES` clause for `LOAD DATA INFILE`. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

- `--quote-names`, `-Q`

  Quote identifiers (such as database, table, and column names) within ` characters. If the `ANSI_QUOTES` SQL mode is enabled, identifiers are quoted within " characters. This option is enabled by default. It can be disabled with `--skip-quote-names`, but this option should be given after any option such as `--compatible` that may enable `--quote-names`.

- `--result-file=`*`file_name`*, `-r` *`file_name`*

  Direct output to the named file. The result file is created and its previous contents overwritten, even if an error occurs while generating the dump.

  This option should be used on Windows to prevent newline `\n` characters from being converted to `\r\n` carriage return/newline sequences.

- `--tab=`*`dir_name`*, `-T` *`dir_name`*

  Produce tab-separated text-format data files. For each dumped table, `mysqldump` creates a *`tbl_name`*`.sql` file that contains the `CREATE TABLE` statement that creates the table, and the server writes a *`tbl_name`*`.txt` file that contains its data. The option value is the directory in which to write the files.

  Note

  This option should be used only when `mysqldump` is run on the same machine as the `mysqld` server. Because the server creates `*.txt` files in the directory that you specify, the directory must be writable by the server and the MySQL account that you use must have the `FILE` privilege. Because `mysqldump` creates `*.sql` in the same directory, it must be writable by your system login account.

  By default, the `.txt` data files are formatted using tab characters between column values and a newline at the end of each line. The format can be specified explicitly using the `--fields-`*`xxx`* and `--lines-terminated-by` options.

  Column values are converted to the character set specified by the `--default-character-set` option.

- `--tz-utc`

  This option enables `TIMESTAMP` columns to be dumped and reloaded between servers in different time zones. `mysqldump` sets its connection time zone to UTC and adds `SET TIME_ZONE='+00:00'` to the dump file. Without this option, `TIMESTAMP` columns are dumped and reloaded in the time zones local to the source and destination servers, which can cause the values to change if the servers are in different time zones. `--tz-utc` also protects against changes due to daylight saving time. `--tz-utc` is enabled by default. To disable it, use `--skip-tz-utc`.

- `--xml`, `-X`

  Write dump output as well-formed XML.

  `NULL`, `'NULL'`, and Empty Values: For a column named *`column_name`*, the `NULL` value, an empty string, and the string value `'NULL'` are distinguished from one another in the output generated by this option as follows.

Value:	XML Representation:
`NULL` (*unknown value*)	`<field name="column_name" xsi:nil="true" />`
`''` (*empty string*)	`<field name="column_name"></field>`
`'NULL'` (*string value*)	`<field name="column_name">NULL</field>`

  The output from the `mysql` client when run using the `--xml` option also follows the preceding rules. (See Section 4.5.1.1, “mysql Options”.)

  XML output from `mysqldump` includes the XML namespace, as shown here:

```
shell> mysqldump --xml -u root world City
<?xml version="1.0"?>
<mysqldump xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance">
<database name="world">
<table_structure name="City">
<field Field="ID" Type="int(11)" Null="NO" Key="PRI" Extra="auto_increment" />
<field Field="Name" Type="char(35)" Null="NO" Key="" Default="" Extra="" />
<field Field="CountryCode" Type="char(3)" Null="NO" Key="" Default="" Extra="" />
<field Field="District" Type="char(20)" Null="NO" Key="" Default="" Extra="" />
<field Field="Population" Type="int(11)" Null="NO" Key="" Default="0" Extra="" />
<key Table="City" Non_unique="0" Key_name="PRIMARY" Seq_in_index="1" Column_name="ID"
Collation="A" Cardinality="4079" Null="" Index_type="BTREE" Comment="" />
<options Name="City" Engine="MyISAM" Version="10" Row_format="Fixed" Rows="4079"
Avg_row_length="67" Data_length="273293" Max_data_length="18858823439613951"
Index_length="43008" Data_free="0" Auto_increment="4080"
Create_time="2007-03-31 01:47:01" Update_time="2007-03-31 01:47:02"
Collation="latin1_swedish_ci" Create_options="" Comment="" />
</table_structure>
<table_data name="City">
<row>
<field name="ID">1</field>
<field name="Name">Kabul</field>
<field name="CountryCode">AFG</field>
<field name="District">Kabol</field>
<field name="Population">1780000</field>
</row>

...

<row>
<field name="ID">4079</field>
<field name="Name">Rafah</field>
<field name="CountryCode">PSE</field>
<field name="District">Rafah</field>
<field name="Population">92020</field>
```

```
</row>
</table_data>
</database>
</mysqldump>
```

Filtering Options

The following options control which kinds of schema objects are written to the dump file: by category, such as triggers or events; by name, for example, choosing which databases and tables to dump; or even filtering rows from the table data using a `WHERE` clause.

- `--all-databases`, `-A`

  Dump all tables in all databases. This is the same as using the `--databases` option and naming all the databases on the command line.

  Prior to MySQL 8.0, the `--routines` and `--events` options for `mysqldump` and `mysqlpump` were not required to include stored routines and events when using the `--all-databases` option: The dump included the `mysql` system database, and therefore also the `mysql.proc` and `mysql.event` tables containing stored routine and event definitions. As of MySQL 8.0, the `mysql.event` and `mysql.proc` tables are not used. Definitions for the corresponding objects are stored in data dictionary tables, but those tables are not dumped. To include stored routines and events in a dump made using `--all-databases`, use the `--routines` and `--events` options explicitly.

- `--databases`, `-B`

  Dump several databases. Normally, `mysqldump` treats the first name argument on the command line as a database name and following names as table names. With this option, it treats all name arguments as database names. `CREATE DATABASE` and `USE` statements are included in the output before each new database.

  This option may be used to dump the `performace_schema` database, which normally is not dumped even with the `--all-databases` option. (Also use the `--skip-lock-tables` option.)

- `--events`, `-E`

  Include Event Scheduler events for the dumped databases in the output. This option requires the `EVENT` privileges for those databases.

  The output generated by using `--events` contains `CREATE EVENT` statements to create the events.

- `--ignore-error=`*`error[,error]...`*

  Ignore the specified errors. The option value is a comma-separated list of error numbers specifying the errors to ignore during `mysqldump` execution. If the `--force` option is also given to ignore all errors, `--force` takes precedence.

- `--ignore-table=`*`db_name.tbl_name`*

  Do not dump the given table, which must be specified using both the database and table names. To ignore multiple tables, use this option multiple times. This option also can be used to ignore views.

- `--no-data`, `-d`

  Do not write any table row information (that is, do not dump table contents). This is useful if you want to dump only the `CREATE TABLE` statement for the table (for example, to create an empty copy of the table by loading the dump file).

- `--routines`, `-R`

  Include stored routines (procedures and functions) for the dumped databases in the output. This option requires the global `SELECT` privilege.

The output generated by using --routines contains CREATE PROCEDURE and CREATE FUNCTION statements to create the routines.

- --tables

  Override the --databases or -B option. mysqldump regards all name arguments following the option as table names.

- --triggers

  Include triggers for each dumped table in the output. This option is enabled by default; disable it with --skip-triggers.

  To be able to dump a table's triggers, you must have the TRIGGER privilege for the table.

  Multiple triggers are permitted. mysqldump dumps triggers in activation order so that when the dump file is reloaded, triggers are created in the same activation order. However, if a mysqldump dump file contains multiple triggers for a table that have the same trigger event and action time, an error occurs for attempts to load the dump file into an older server that does not support multiple triggers. (For a workaround, see Section 2.10.2.1, “Changes Affecting Downgrades from MySQL 8.0”; you can convert triggers to be compatible with older servers.)

- --where='*where_condition*', -w '*where_condition*'

  Dump only rows selected by the given WHERE condition. Quotes around the condition are mandatory if it contains spaces or other characters that are special to your command interpreter.

  Examples:

```
--where="user='jimf'"
-w"userid>1"
-w"userid<1"
```

Performance Options

The following options are the most relevant for the performance particularly of the restore operations. For large data sets, restore operation (processing the INSERT statements in the dump file) is the most time-consuming part. When it is urgent to restore data quickly, plan and test the performance of this stage in advance. For restore times measured in hours, you might prefer an alternative backup and restore solution, such as MySQL Enterprise Backup for InnoDB-only and mixed-use databases.

Performance is also affected by the transactional options, primarily for the dump operation.

- --column-statistics

  Add ANALYZE TABLE statements to the output to generate histogram statistics for dumped tables when the dump file is reloaded. This option is disabled by default because histogram generation for large tables can take a long time.

  This option was added in MySQL 8.0.2.

- --disable-keys, -K

  For each table, surround the INSERT statements with /*!40000 ALTER TABLE *tbl_name* DISABLE KEYS */; and /*!40000 ALTER TABLE *tbl_name* ENABLE KEYS */; statements. This makes loading the dump file faster because the indexes are created after all rows are inserted. This option is effective only for nonunique indexes of MyISAM tables.

- --extended-insert, -e

Write INSERT statements using multiple-row syntax that includes several VALUES lists. This results in a smaller dump file and speeds up inserts when the file is reloaded.

- --insert-ignore

  Write INSERT IGNORE statements rather than INSERT statements.

- --opt

  This option, enabled by default, is shorthand for the combination of --add-drop-table --add-locks --create-options --disable-keys --extended-insert --lock-tables --quick --set-charset. It gives a fast dump operation and produces a dump file that can be reloaded into a MySQL server quickly.

  Because the --opt option is enabled by default, you only specify its converse, the --skip-opt to turn off several default settings. See the discussion of mysqldump option groups for information about selectively enabling or disabling a subset of the options affected by --opt.

- --quick, -q

  This option is useful for dumping large tables. It forces mysqldump to retrieve rows for a table from the server a row at a time rather than retrieving the entire row set and buffering it in memory before writing it out.

- --skip-opt

  See the description for the --opt option.

Transactional Options

The following options trade off the performance of the dump operation, against the reliability and consistency of the exported data.

- --add-locks

  Surround each table dump with LOCK TABLES and UNLOCK TABLES statements. This results in faster inserts when the dump file is reloaded. See Section 8.2.5.1, “Optimizing INSERT Statements”.

- --flush-logs, -F

  Flush the MySQL server log files before starting the dump. This option requires the RELOAD privilege. If you use this option in combination with the --all-databases option, the logs are flushed *for each database dumped*. The exception is when using --lock-all-tables, --master-data, or --single-transaction: In this case, the logs are flushed only once, corresponding to the moment that all tables are locked by FLUSH TABLES WITH READ LOCK. If you want your dump and the log flush to happen at exactly the same moment, you should use --flush-logs together with --lock-all-tables, --master-data, or --single-transaction.

- --flush-privileges

  Add a FLUSH PRIVILEGES statement to the dump output after dumping the mysql database. This option should be used any time the dump contains the mysql database and any other database that depends on the data in the mysql database for proper restoration.

  Note

  For upgrades to MySQL 5.7.2 or higher from older versions, do not use --flush-privileges. For upgrade instructions in this case, see Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0”.

- --lock-all-tables, -x

Lock all tables across all databases. This is achieved by acquiring a global read lock for the duration of the whole dump. This option automatically turns off `--single-transaction` and `--lock-tables`.

- `--lock-tables`, `-l`

  For each dumped database, lock all tables to be dumped before dumping them. The tables are locked with `READ LOCAL` to permit concurrent inserts in the case of `MyISAM` tables. For transactional tables such as `InnoDB`, `--single-transaction` is a much better option than `--lock-tables` because it does not need to lock the tables at all.

  Because `--lock-tables` locks tables for each database separately, this option does not guarantee that the tables in the dump file are logically consistent between databases. Tables in different databases may be dumped in completely different states.

  Some options, such as `--opt`, automatically enable `--lock-tables`. If you want to override this, use `--skip-lock-tables` at the end of the option list.

- `--no-autocommit`

  Enclose the `INSERT` statements for each dumped table within `SET autocommit = 0` and `COMMIT` statements.

- `--order-by-primary`

  Dump each table's rows sorted by its primary key, or by its first unique index, if such an index exists. This is useful when dumping a `MyISAM` table to be loaded into an `InnoDB` table, but makes the dump operation take considerably longer.

- `--shared-memory-base-name=`*name*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--single-transaction`

  This option sets the transaction isolation mode to `REPEATABLE READ` and sends a `START TRANSACTION` SQL statement to the server before dumping data. It is useful only with transactional tables such as `InnoDB`, because then it dumps the consistent state of the database at the time when `START TRANSACTION` was issued without blocking any applications.

  When using this option, you should keep in mind that only `InnoDB` tables are dumped in a consistent state. For example, any `MyISAM` or `MEMORY` tables dumped while using this option may still change state.

  While a `--single-transaction` dump is in process, to ensure a valid dump file (correct table contents and binary log coordinates), no other connection should use the following statements: `ALTER TABLE`, `CREATE TABLE`, `DROP TABLE`, `RENAME TABLE`, `TRUNCATE TABLE`. A consistent read is not isolated from those statements, so use of them on a table to be dumped can cause the `SELECT` that is performed by `mysqldump` to retrieve the table contents to obtain incorrect contents or fail.

  The `--single-transaction` option and the `--lock-tables` option are mutually exclusive because `LOCK TABLES` causes any pending transactions to be committed implicitly.

  To dump large tables, combine the `--single-transaction` option with the `--quick` option.

Option Groups

- The `--opt` option turns on several settings that work together to perform a fast dump operation. All of these settings are on by default, because `--opt` is on by default. Thus you rarely if ever specify `--opt`. Instead, you can turn these settings off as a group by specifying `--skip-opt`, the optionally re-enable certain settings by specifying the associated options later on the command line.
- The `--compact` option turns off several settings that control whether optional statements and comments appear in the output. Again, you can follow this option with other options that re-enable certain settings, or turn all the settings on by using the `--skip-compact` form.

When you selectively enable or disable the effect of a group option, order is important because options are processed first to last. For example, `--disable-keys --lock-tables --skip-opt` would not have the intended effect; it is the same as `--skip-opt` by itself.

Examples

To make a backup of an entire database:

```
shell> mysqldump db_name > backup-file.sql
```

To load the dump file back into the server:

```
shell> mysql db_name < backup-file.sql
```

Another way to reload the dump file:

```
shell> mysql -e "source /path-to-backup/backup-file.sql" db_name
```

`mysqldump` is also very useful for populating databases by copying data from one MySQL server to another:

```
shell> mysqldump --opt db_name | mysql --host=remote_host -C db_name
```

You can dump several databases with one command:

```
shell> mysqldump --databases db_name1 [db_name2 ...] > my_databases.sql
```

To dump all databases, use the `--all-databases` option:

```
shell> mysqldump --all-databases > all_databases.sql
```

For `InnoDB` tables, `mysqldump` provides a way of making an online backup:

```
shell> mysqldump --all-databases --master-data --single-transaction > all_databases.sql
```

This backup acquires a global read lock on all tables (using `FLUSH TABLES WITH READ LOCK`) at the beginning of the dump. As soon as this lock has been acquired, the binary log coordinates are read and the lock is released. If long updating statements are running when the `FLUSH` statement is issued, the MySQL server may get stalled until those statements finish. After that, the dump becomes lock free and does not disturb reads and writes on the tables. If the update statements that the MySQL server receives are short (in terms of execution time), the initial lock period should not be noticeable, even with many updates.

For point-in-time recovery (also known as “roll-forward,” when you need to restore an old backup and replay the changes that happened since that backup), it is often useful to rotate the binary log

(see Section 5.4.4, “The Binary Log”) or at least know the binary log coordinates to which the dump corresponds:

```
shell> mysqldump --all-databases --master-data=2 > all_databases.sql
```

Or:

```
shell> mysqldump --all-databases --flush-logs --master-data=2
              > all_databases.sql
```

The `--master-data` and `--single-transaction` options can be used simultaneously, which provides a convenient way to make an online backup suitable for use prior to point-in-time recovery if tables are stored using the `InnoDB` storage engine.

For more information on making backups, see Section 7.2, “Database Backup Methods”, and Section 7.3, “Example Backup and Recovery Strategy”.

- To select the effect of `--opt` except for some features, use the `--skip` option for each feature. To disable extended inserts and memory buffering, use `--opt --skip-extended-insert --skip-quick`. (Actually, `--skip-extended-insert --skip-quick` is sufficient because `--opt` is on by default.)
- To reverse `--opt` for all features except index disabling and table locking, use `--skip-opt --disable-keys --lock-tables`.

Restrictions

`mysqldump` does not dump the `performance_schema` or `sys` schema by default. To dump any of these, name them explicitly on the command line. You can also name them with the `--databases` option. For `performance_schema`, also use the `--skip-lock-tables` option.

`mysqldump` does not dump the `INFORMATION_SCHEMA` schema.

`mysqldump` includes statements to recreate the `general_log` and `slow_query_log` tables for dumps of the `mysql` database. Log table contents are not dumped.

If you encounter problems backing up views due to insufficient privileges, see Section C.5, “Restrictions on Views” for a workaround.

4.5.5 `mysqlimport` — A Data Import Program

The `mysqlimport` client provides a command-line interface to the `LOAD DATA INFILE` SQL statement. Most options to `mysqlimport` correspond directly to clauses of `LOAD DATA INFILE` syntax. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

Invoke `mysqlimport` like this:

```
shell> mysqlimport [options] db_name textfile1 [textfile2 ...]
```

For each text file named on the command line, `mysqlimport` strips any extension from the file name and uses the result to determine the name of the table into which to import the file's contents. For example, files named `patient.txt`, `patient.text`, and `patient` all would be imported into a table named `patient`.

`mysqlimport` supports the following options, which can be specified on the command line or in the `[mysqlimport]` and `[client]` groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.12 mysqlimport Options

Format	Description
--bind-address	Use specified network interface to connect to MySQL Server
--columns	This option takes a comma-separated list of column names as its value
--compress	Compress all information sent between client and server
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--default-character-set	Specify default character set
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--delete	Empty the table before importing the text file
--enable-cleartext-plugin	Enable cleartext authentication plugin
--fields-enclosed-by	This option has the same meaning as the corresponding clause for LOAD DATA INFILE
--fields-escaped-by	This option has the same meaning as the corresponding clause for LOAD DATA INFILE
--fields-optionally-enclosed-by	This option has the same meaning as the corresponding clause for LOAD DATA INFILE
--fields-terminated-by	This option has the same meaning as the corresponding clause for LOAD DATA INFILE
--force	Continue even if an SQL error occurs
--help	Display help message and exit
--host	Connect to MySQL server on given host
--ignore	See the description for the --replace option
--ignore-lines	Ignore the first N lines of the data file
--lines-terminated-by	This option has the same meaning as the corresponding clause for LOAD DATA INFILE
--local	Read input files locally from the client host
--lock-tables	Lock all tables for writing before processing any text files
--login-path	Read login path options from .mylogin.cnf
--low-priority	Use LOW_PRIORITY when loading the table.
--no-defaults	Read no option files
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use

Format	Description
--replace	The --replace and --ignore options control handling of input rows that duplicate existing rows on unique key values
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--silent	Produce output only when errors occur
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tls-version	Protocols permitted for secure connections
--use-threads	Number of threads for parallel file-loading
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version	Display version information and exit

- `--help`, `-?`

  Display a help message and exit.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--columns=`*`column_list`*, `-c` *`column_list`*

  This option takes a comma-separated list of column names as its value. The order of the column names indicates how to match data file columns with table columns.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--debug[=`*`debug_options`*`]`, `-#` `[`*`debug_options`*`]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:o`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, “Character Set Configuration”.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqlimport` normally reads the `[client]` and `[mysqlimport]` groups. If the `--defaults-group-suffix=_other` option is given, `mysqlimport` also reads the `[client_other]` and `[mysqlimport_other]` groups.

- `--delete`, `-D`

  Empty the table before importing the text file.

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--fields-terminated-by=...`, `--fields-enclosed-by=...`, `--fields-optionally-enclosed-by=...`, `--fields-escaped-by=...`

  These options have the same meaning as the corresponding clauses for `LOAD DATA INFILE`. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

- `--force`, `-f`

  Ignore errors. For example, if a table for a text file does not exist, continue processing any remaining files. Without `--force`, `mysqlimport` exits if a table does not exist.

- `--host=host_name`, `-h host_name`

  Import data to the MySQL server on the given host. The default host is `localhost`.

- `--ignore`, `-i`

  See the description for the `--replace` option.

- `--ignore-lines=N`

Ignore the first *N* lines of the data file.

- `--lines-terminated-by=...`

  This option has the same meaning as the corresponding clause for `LOAD DATA INFILE`. For example, to import Windows files that have lines terminated with carriage return/linefeed pairs, use `--lines-terminated-by="\r\n"`. (You might have to double the backslashes, depending on the escaping conventions of your command interpreter.) See Section 13.2.7, “LOAD DATA INFILE Syntax”.

- `--local`, `-L`

  By default, files are read by the server on the server host. With this option, `mysqlimport` reads input files locally on the client host. Enabling local data loading also requires that the server permits it; see Section 6.1.6, “Security Issues with LOAD DATA LOCAL”

- `--lock-tables`, `-l`

  Lock *all* tables for writing before processing any text files. This ensures that all tables are synchronized on the server.

- `--login-path=`*`name`*

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--low-priority`

  Use `LOW_PRIORITY` when loading the table. This affects only storage engines that use only table-level locking (such as `MyISAM`, `MEMORY`, and `MERGE`).

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.)

- `--password[=`*`password`*`]`, `-p[`*`password`*`]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysqlimport` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=`*`dir_name`*

The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqlimport` does not find it. See Section 6.3.10, “Pluggable Authentication”.

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--replace`, `-r`

  The `--replace` and `--ignore` options control handling of input rows that duplicate existing rows on unique key values. If you specify `--replace`, new rows replace existing rows that have the same unique key value. If you specify `--ignore`, input rows that duplicate an existing row on a unique key value are skipped. If you do not specify either option, an error occurs when a duplicate key value is found, and the rest of the text file is ignored.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--shared-memory-base-name=name`

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--silent`, `-s`

  Silent mode. Produce output only when errors occur.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=protocol_list`

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server.

- `--use-threads=N`

  Load files in parallel using `N` threads.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does.

- `--version`, `-V`

  Display version information and exit.

Here is a sample session that demonstrates use of `mysqlimport`:

```
shell> mysql -e 'CREATE TABLE imptest(id INT, n VARCHAR(30))' test
shell> ed
a
100     Max Sydow
101     Count Dracula
.
w imptest.txt
32
q
shell> od -c imptest.txt
0000000   1   0   0  \t   M   a   x       S   y   d   o   w  \n   1   0
0000020   1  \t   C   o   u   n   t       D   r   a   c   u   l   a  \n
0000040
shell> mysqlimport --local test imptest.txt
test.imptest: Records: 2  Deleted: 0  Skipped: 0  Warnings: 0
shell> mysql -e 'SELECT * FROM imptest' test
+------+---------------+
| id   | n             |
+------+---------------+
|  100 | Max Sydow     |
|  101 | Count Dracula |
+------+---------------+
```

4.5.6 mysqlpump — A Database Backup Program

- mysqlpump Invocation Syntax
- mysqlpump Option Summary
- mysqlpump Option Descriptions
- mysqlpump Object Selection
- mysqlpump Parallel Processing
- mysqlpump Restrictions

The mysqlpump client utility performs logical backups, producing a set of SQL statements that can be executed to reproduce the original database object definitions and table data. It dumps one or more MySQL databases for backup or transfer to another SQL server.

mysqlpump features include:

- Parallel processing of databases, and of objects within databases, to speed up the dump process
- Better control over which databases and database objects (tables, stored programs, user accounts) to dump
- Dumping of user accounts as account-management statements (CREATE USER, GRANT) rather than as inserts into the mysql system database
- Capability of creating compressed output
- Progress indicator (the values are estimates)
- For dump file reloading, faster secondary index creation for InnoDB tables by adding indexes after rows are inserted

Note

mysqlpump uses MySQL features introduced in MySQL 5.7, and thus assumes use with MySQL 5.7 or higher.

mysqlpump requires at least the SELECT privilege for dumped tables, SHOW VIEW for dumped views, TRIGGER for dumped triggers, and LOCK TABLES if the --single-transaction option is not used. The SELECT privilege on the mysql system database is required to dump user definitions. Certain options might require other privileges as noted in the option descriptions.

To reload a dump file, you must have the privileges required to execute the statements that it contains, such as the appropriate CREATE privileges for objects created by those statements.

Note

A dump made using PowerShell on Windows with output redirection creates a file that has UTF-16 encoding:

```
shell> mysqlpump [options] > dump.sql
```

However, UTF-16 is not permitted as a connection character set (see Section 10.1.4, “Connection Character Sets and Collations”), so the dump file will not load correctly. To work around this issue, use the --result-file option, which creates the output in ASCII format:

```
shell> mysqlpump [options] --result-file=dump.sql
```

mysqlpump Invocation Syntax

By default, mysqlpump dumps all databases (with certain exceptions noted in mysqlpump Restrictions). To specify this behavior explicitly, use the --all-databases option:

```
shell> mysqlpump --all-databases
```

To dump a single database, or certain tables within that database, name the database on the command line, optionally followed by table names:

```
shell> mysqlpump db_name
shell> mysqlpump db_name tbl_name1 tbl_name2 ...
```

To treat all name arguments as database names, use the --databases option:

```
shell> mysqlpump --databases db_name1 db_name2 ...
```

By default, mysqlpump does not dump user account definitions, even if you dump the mysql system database that contains the grant tables. To dump grant table contents as logical definitions in the form of CREATE USER and GRANT statements, use the --users option and suppress all database dumping:

```
shell> mysqlpump --exclude-databases=% --users
```

In the preceding command, % is a wildcard that matches all database names for the --exclude-databases option.

mysqlpump supports several options for including or excluding databases, tables, stored programs, and user definitions. See mysqlpump Object Selection.

To reload a dump file, execute the statements that it contains. For example, use the mysql client:

```
shell> mysqlpump [options] > dump.sql
shell> mysql < dump.sql
```

The following discussion provides additional mysqlpump usage examples.

To see a list of the options mysqlpump supports, issue the command mysqlpump --help.

mysqlpump Option Summary

mysqlpump supports the following options, which can be specified on the command line or in the [mysqlpump] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.13 mysqlpump Options

Format	Description	Introduced
--add-drop-database	Add DROP DATABASE statement before each CREATE DATABASE statement	
--add-drop-table	Add DROP TABLE statement before each CREATE TABLE statement	
--add-drop-user	Add DROP USER statement before each CREATE USER statement	
--add-locks	Surround each table dump with LOCK TABLES and UNLOCK TABLES statements	
--all-databases	Dump all databases	
--bind-address	Use specified network interface to connect to MySQL Server	
--character-sets-dir	Directory where character sets are installed	
--column-statistics	Write ANALYZE TABLE statements to generate statistics histograms	8.0.2
--complete-insert	Use complete INSERT statements that include column names	
--compress	Compress all information sent between client and server	
--compress-output	Output compression algorithm	
--databases	Interpret all name arguments as database names	
--debug	Write debugging log	
--debug-check	Print debugging information when program exits	

Format	Description	Introduced
--debug-info	Print debugging information, memory, and CPU statistics when program exits	
--default-auth	Authentication plugin to use	
--default-character-set	Specify default character set	
--default-parallelism	Default number of threads for parallel processing	
--defaults-extra-file	Read named option file in addition to usual option files	
--defaults-file	Read only named option file	
--defaults-group-suffix	Option group suffix value	
--defer-table-indexes	For reloading, defer index creation until after loading table rows	
--events	Dump events from dumped databases	
--exclude-databases	Databases to exclude from dump	
--exclude-events	Events to exclude from dump	
--exclude-routines	Routines to exclude from dump	
--exclude-tables	Tables to exclude from dump	
--exclude-triggers	Triggers to exclude from dump	
--exclude-users	Users to exclude from dump	
--extended-insert	Use multiple-row INSERT syntax	
--help	Display help message and exit	
--hex-blob	Dump binary columns using hexadecimal notation	
--host	Host to connect to (IP address or hostname)	
--include-databases	Databases to include in dump	
--include-events	Events to include in dump	
--include-routines	Routines to include in dump	
--include-tables	Tables to include in dump	
--include-triggers	Triggers to include in dump	
--include-users	Users to include in dump	
--insert-ignore	Write INSERT IGNORE rather than INSERT statements	
--log-error-file	Append warnings and errors to named file	
--login-path	Read login path options from .mylogin.cnf	
--max-allowed-packet	Maximum packet length to send to or receive from server	
--net-buffer-length	Buffer size for TCP/IP and socket communication	
--no-create-db	Do not write CREATE DATABASE statements	
--no-create-info	Do not write CREATE TABLE statements that re-create each dumped table	
--no-defaults	Read no option files	
--parallel-schemas	Specify schema-processing parallelism	
--password	Password to use when connecting to server	
--plugin-dir	Directory where plugins are installed	
--port	TCP/IP port number to use for connection	
--print-defaults	Print default options	
--protocol	Connection protocol to use	

Format	Description	Introduced
--replace	Write REPLACE statements rather than INSERT statements	
--result-file	Direct output to a given file	
--routines	Dump stored routines (procedures and functions) from dumped databases	
--secure-auth	Do not send passwords to server in old (pre-4.1) format	
--set-charset	Add SET NAMES default_character_set to output	
--set-gtid-purged	Whether to add SET @@GLOBAL.GTID_PURGED to output	8.0.1
--single-transaction	Dump tables within single transaction	
--skip-definer	Omit DEFINER and SQL SECURITY clauses from view and stored program CREATE statements	
--skip-dump-rows	Do not dump table rows	
--socket	For connections to localhost, the Unix socket file to use	
--ssl-ca	Path of file that contains list of trusted SSL CAs	
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format	
--ssl-cert	Path of file that contains X509 certificate in PEM format	
--ssl-cipher	List of permitted ciphers to use for connection encryption	
--ssl-crl	Path of file that contains certificate revocation lists	
--ssl-crlpath	Path of directory that contains certificate revocation list files	
--ssl-key	Path of file that contains X509 key in PEM format	
--ssl-mode	Security state of connection to server	
--tls-version	Protocols permitted for secure connections	
--triggers	Dump triggers for each dumped table	
--tz-utc	Add SET TIME_ZONE='+00:00' to dump file	
--user	MySQL user name to use when connecting to server	
--users	Dump user accounts	
--version	Display version information and exit	
--watch-progress	Display progress indicator	

mysqlpump Option Descriptions

- `--help`, `-?`

  Display a help message and exit.

- `--add-drop-database`

  Write a `DROP DATABASE` statement before each `CREATE DATABASE` statement.

- `--add-drop-table`

  Write a `DROP TABLE` statement before each `CREATE TABLE` statement.

- `--add-drop-user`

  Write a `DROP USER` statement before each `CREATE USER` statement.

- `--add-locks`

  Surround each table dump with `LOCK TABLES` and `UNLOCK TABLES` statements. This results in faster inserts when the dump file is reloaded. See Section 8.2.5.1, “Optimizing INSERT Statements”.

  This option does not work with parallelism because `INSERT` statements from different tables can be interleaved and `UNLOCK TABLES` following the end of the inserts for one table could release locks on tables for which inserts remain.

  `--add-locks` and `--single-transaction` are mutually exclusive.

- `--all-databases`, `-A`

  Dump all databases (with certain exceptions noted in mysqlpump Restrictions). This is the default behavior if no other is specified explicitly.

  `--all-databases` and `--databases` are mutually exclusive.

  Prior to MySQL 8.0, the `--routines` and `--events` options for `mysqldump` and `mysqlpump` were not required to include stored routines and events when using the `--all-databases` option: The dump included the `mysql` system database, and therefore also the `mysql.proc` and `mysql.event` tables containing stored routine and event definitions. As of MySQL 8.0, the `mysql.event` and `mysql.proc` tables are not used. Definitions for the corresponding objects are stored in data dictionary tables, but those tables are not dumped. To include stored routines and events in a dump made using `--all-databases`, use the `--routines` and `--events` options explicitly.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`path`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--column-statistics`

  Add `ANALYZE TABLE` statements to the output to generate histogram statistics for dumped tables when the dump file is reloaded. This option is disabled by default because histogram generation for large tables can take a long time.

  This option was added in MySQL 8.0.2.

- `--complete-insert`

  Write complete `INSERT` statements that include column names.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--compress-output=`*`algorithm`*

  By default, `mysqlpump` does not compress output. This option specifies output compression using the specified algorithm. Permitted algorithms are `LZ4` and `ZLIB`.

  To uncompress compressed output, you must have an appropriate utility. If the system commands `lz4` and `openssl zlib` are not available, MySQL distributions include `lz4_decompress` and `zlib_decompress` utilities that can be used to decompress `mysqlpump` output that was compressed using the `--compress-output=LZ4` and `--compress-output=ZLIB` options.

  For more information, see Section 4.8.1, “`lz4_decompress` — Decompress mysqlpump LZ4 Output”.

Compressed Output", and Section 4.8.4, "zlib_decompress — Decompress mysqlpump ZLIB-Compressed Output".

- `--databases`, `-B`

  Normally, `mysqlpump` treats the first name argument on the command line as a database name and any following names as table names. With this option, it treats all name arguments as database names. `CREATE DATABASE` statements are included in the output before each new database.

  `--all-databases` and `--databases` are mutually exclusive.

- `--debug[=debug_options]`, `-# [debug_options]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:O,/tmp/mysqlpump.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`, `-T`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, "Pluggable Authentication".

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, "Character Set Configuration". If no character set is specified, `mysqlpump` uses `utf8`.

- `--default-parallelism=N`

  The default number of threads for each parallel processing queue. The default is 2.

  The `--parallel-schemas` option also affects parallelism and can be used to override the default number of threads. For more information, see mysqlpump Parallel Processing.

  With `--default-parallelism=0` and no `--parallel-schemas` options, `mysqlpump` runs as a single-threaded process and creates no queues.

  With parallelism enabled, it is possible for output from different databases to be interleaved.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

Read not only the usual option groups, but also groups with the usual names and a suffix of *str*. For example, mysqlpump normally reads the [client] and [mysqlpump] groups. If the --defaults-group-suffix=_other option is given, mysqlpump also reads the [client_other] and [mysqlpump_other] groups.

- --defer-table-indexes

  In the dump output, defer index creation for each table until after its rows have been loaded. This works for all storage engines, but for InnoDB applies only for secondary indexes.

  This option is enabled by default; use --skip-defer-table-indexes to disable it.

- --events

  Include Event Scheduler events for the dumped databases in the output. Event dumping requires the EVENT privileges for those databases.

  The output generated by using --events contains CREATE EVENT statements to create the events.

  This option is enabled by default; use --skip-events to disable it.

- --exclude-databases=*db_list*

  Do not dump the databases in *db_list*, which is a comma-separated list of one or more database names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --exclude-events=*event_list*

  Do not dump the databases in *event_list*, which is a comma-separated list of one or more event names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --exclude-routines=*routine_list*

  Do not dump the events in *routine_list*, which is a comma-separated list of one or more routine (stored procedure or function) names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --exclude-tables=*table_list*

  Do not dump the tables in *table_list*, which is a comma-separated list of one or more table names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --exclude-triggers=*trigger_list*

  Do not dump the triggers in *trigger_list*, which is a comma-separated list of one or more trigger names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --exclude-users=*user_list*

  Do not dump the user accounts in *user_list*, which is a comma-separated list of one or more account names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --extended-insert=*N*

  Write INSERT statements using multiple-row syntax that includes several VALUES lists. This results in a smaller dump file and speeds up inserts when the file is reloaded.

The option value indicates the number of rows to include in each INSERT statement. The default is 250. A value of 1 produces one INSERT statement per table row.

- --hex-blob

  Dump binary columns using hexadecimal notation (for example, 'abc' becomes 0x616263). The affected data types are BINARY, VARBINARY, the BLOB types, and BIT.

- --host=*host_name*, -h *host_name*

  Dump data from the MySQL server on the given host.

- --include-databases=*db_list*

  Dump the databases in *db_list*, which is a comma-separated list of one or more database names. The dump includes all objects in the named databases. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --include-events=*event_list*

  Dump the events in *event_list*, which is a comma-separated list of one or more event names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --include-routines=*routine_list*

  Dump the routines in *routine_list*, which is a comma-separated list of one or more routine (stored procedure or function) names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --include-tables=*table_list*

  Dump the tables in *table_list*, which is a comma-separated list of one or more table names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --include-triggers=*trigger_list*

  Dump the triggers in *trigger_list*, which is a comma-separated list of one or more trigger names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --include-users=*user_list*

  Dump the user accounts in *user_list*, which is a comma-separated list of one or more user names. Multiple instances of this option are additive. For more information, see mysqlpump Object Selection.

- --insert-ignore

  Write INSERT IGNORE statements rather than INSERT statements.

- --log-error-file=*file_name*

  Log warnings and errors by appending them to the named file. If this option is not given, mysqlpump writes warnings and errors to the standard error output.

- --login-path=*name*

  Read options from the named login path in the .mylogin.cnf login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the mysql_config_editor utility. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

- --max-allowed-packet=*N*

The maximum size of the buffer for client/server communication. The default is 24MB, the maximum is 1GB.

- `--net-buffer-length=N`

  The initial size of the buffer for client/server communication. When creating multiple-row `INSERT` statements (as with the `--extended-insert` option), `mysqlpump` creates rows up to *`N`* bytes long. If you use this option to increase the value, ensure that the MySQL server `net_buffer_length` system variable has a value at least this large.

- `--no-create-db`

  Suppress any `CREATE DATABASE` statements that might otherwise be included in the output.

- `--no-create-info`, `-t`

  Do not write `CREATE TABLE` statements that create each dumped table.

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.)

- `--parallel-schemas=[N:]db_list`

  Create a queue for processing the databases in *`db_list`*, which is a comma-separated list of one or more database names. If *`N`* is given, the queue uses *`N`* threads. If *`N`* is not given, the `--default-parallelism` option determines the number of queue threads.

  Multiple instances of this option create multiple queues. `mysqlpump` also creates a default queue to use for databases not named in any `--parallel-schemas` option, and for dumping user definitions if command options select them. For more information, see mysqlpump Parallel Processing.

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysqlpump` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file to avoid giving the password on the command line.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqlpump` does not find it. See Section 6.3.10, “Pluggable Authentication”.

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--replace`

  Write `REPLACE` statements rather than `INSERT` statements.

- `--result-file=file_name`

  Direct output to the named file. The result file is created and its previous contents overwritten, even if an error occurs while generating the dump.

  This option should be used on Windows to prevent newline `\n` characters from being converted to `\r\n` carriage return/newline sequences.

- `--routines`

  Include stored routines (procedures and functions) for the dumped databases in the output. This option requires the global `SELECT` privilege.

  The output generated by using `--routines` contains `CREATE PROCEDURE` and `CREATE FUNCTION` statements to create the routines.

  This option is enabled by default; use `--skip-routines` to disable it.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

- `--set-charset`

  Write `SET NAMES default_character_set` to the output.

  This option is enabled by default. To disable it and suppress the `SET NAMES` statement, use `--skip-set-charset`.

- `--set-gtid-purged=value`

  This option enables control over global transaction ID (GTID) information written to the dump file, by indicating whether to add a `SET @@global.gtid_purged` statement to the output. This option may also cause a statement to be written to the output that disables binary logging while the dump file is being reloaded.

  The following table shows the permitted option values. The default value is `AUTO`.

Value	Meaning
`OFF`	Add no `SET` statement to the output.
`ON`	Add a `SET` statement to the output. An error occurs if GTIDs are not enabled on the server.
`AUTO`	Add a `SET` statement to the output if GTIDs are enabled on the server.

  The `--set-gtid-purged` option has the following effect on binary logging when the dump file is reloaded:

  - `--set-gtid-purged=OFF`: `SET @@SESSION.SQL_LOG_BIN=0;` is not added to the output.
  - `--set-gtid-purged=ON`: `SET @@SESSION.SQL_LOG_BIN=0;` is added to the output.
  - `--set-gtid-purged=AUTO`: `SET @@SESSION.SQL_LOG_BIN=0;` is added to the output if GTIDs are enabled on the server you are backing up (that is, if `AUTO` evaluates to `ON`).

- `--single-transaction`

  This option sets the transaction isolation mode to `REPEATABLE READ` and sends a `START TRANSACTION` SQL statement to the server before dumping data. It is useful only with transactional tables such as `InnoDB`, because then it dumps the consistent state of the database at the time when `START TRANSACTION` was issued without blocking any applications.

  When using this option, you should keep in mind that only `InnoDB` tables are dumped in a consistent state. For example, any `MyISAM` or `MEMORY` tables dumped while using this option may still change state.

  While a `--single-transaction` dump is in process, to ensure a valid dump file (correct table contents and binary log coordinates), no other connection should use the following statements: `ALTER TABLE`, `CREATE TABLE`, `DROP TABLE`, `RENAME TABLE`, `TRUNCATE TABLE`. A consistent read is not isolated from those statements, so use of them on a table to be dumped can cause the `SELECT` that is performed by `mysqlpump` to retrieve the table contents to obtain incorrect contents or fail.

  `--add-locks` and `--single-transaction` are mutually exclusive.

- `--skip-definer`

  Omit `DEFINER` and `SQL SECURITY` clauses from the `CREATE` statements for views and stored programs. The dump file, when reloaded, creates objects that use the default `DEFINER` and `SQL SECURITY` values. See Section 23.6, “Access Control for Stored Programs and Views”.

- `--skip-dump-rows`, `-d`

  Do not dump table rows.

- `--socket={`*`file_name`*`|`*`pipe_name`*`}`, `-S {`*`file_name`*`|`*`pipe_name`*`}`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=`*`protocol_list`*

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--triggers`

  Include triggers for each dumped table in the output.

  This option is enabled by default; use `--skip-triggers` to disable it.

- `--tz-utc`

This option enables TIMESTAMP columns to be dumped and reloaded between servers in different time zones. mysqlpump sets its connection time zone to UTC and adds SET TIME_ZONE='+00:00' to the dump file. Without this option, TIMESTAMP columns are dumped and reloaded in the time zones local to the source and destination servers, which can cause the values to change if the servers are in different time zones. --tz-utc also protects against changes due to daylight saving time.

This option is enabled by default; use --skip-tz-utc to disable it.

- --user=*user_name*, -u *user_name*

  The MySQL user name to use when connecting to the server.

- --users

  Dump user accounts as logical definitions in the form of CREATE USER and GRANT statements.

  User definitions are stored in the grant tables in the mysql system database. By default, mysqlpump does not include the grant tables in mysql database dumps. To dump the contents of the grant tables as logical definitions, use the --users option and suppress all database dumping:

```
shell> mysqlpump --exclude-databases=% --users
```

- --version, -V

  Display version information and exit.

- --watch-progress

  Periodically display a progress indicator that provides information about the completed and total number of tables, rows, and other objects.

  This option is enabled by default; use --skip-watch-progress to disable it.

mysqlpump Object Selection

mysqlpump has a set of inclusion and exclusion options that enable filtering of several object types and provide flexible control over which objects to dump:

- --include-databases and --exclude-databases apply to databases and all objects within them.
- --include-tables and --exclude-tables apply to tables. These options also affect triggers associated with tables unless the trigger-specific options are given.
- --include-triggers and --exclude-triggers apply to triggers.
- --include-routines and --exclude-routines apply to stored procedures and functions. If a routine option matches a stored procedure name, it also matches a stored function of the same name.
- --include-events and --exclude-events apply to Event Scheduler events.
- --include-users and --exclude-users apply to user accounts.

Any inclusion or exclusion option may be given multiple times. The effect is additive. Order of these options does not matter.

The value of each inclusion and exclusion option is a comma-separated list of names of the appropriate object type. For example:

```
--exclude-databases=test,world
--include-tables=customer,invoice
```

Wildcard characters are permitted in the object names:

- % matches any sequence of zero or more characters.
- _ matches any single character.

For example, `--include-tables=t%,__tmp` matches all table names that begin with `t` and all five-character table names that end with `tmp`.

For users, a name specified without a host part is interpreted with an implied host of `%`. For example, `u1` and `u1@%` are equivalent. This is the same equivalence that applies in MySQL generally (see Section 6.2.4, “Specifying Account Names”).

Inclusion and exclusion options interact as follows:

- By default, with no inclusion or exclusion options, `mysqlpump` dumps all databases (with certain exceptions noted in mysqlpump Restrictions).
- If inclusion options are given in the absence of exclusion options, only the objects named as included are dumped.
- If exclusion options are given in the absence of inclusion options, all objects are dumped except those named as excluded.
- If inclusion and exclusion options are given, all objects named as excluded and not named as included are not dumped. All other objects are dumped.

If multiple databases are being dumped, it is possible to name tables, triggers, and routines in a specific database by qualifying the object names with the database name. The following command dumps databases `db1` and `db2`, but excludes tables `db1.t1` and `db2.t2`:

```
shell> mysqlpump --include-databases=db1,db2 --exclude-tables=db1.t1,db2.t2
```

The following options provide alternative ways to specify which databases to dump:

- The `--all-databases` option dumps all databases (with certain exceptions noted in mysqlpump Restrictions). It is equivalent to specifying no object options at all (the default `mysqlpump` action is to dump everything).

  `--include-databases=%` is similar to `--all-databases`, but selects all databases for dumping, even those that are exceptions for `--all-databases`.

- The `--databases` option causes `mysqlpump` to treat all name arguments as names of databases to dump. It is equivalent to an `--include-databases` option that names the same databases.

mysqlpump Parallel Processing

`mysqlpump` can use parallelism to achieve concurrent processing. You can select concurrency between databases (to dump multiple databases simultaneously) and within databases (to dump multiple objects from a given database simultaneously).

By default, `mysqlpump` sets up one queue with two threads. You can create additional queues and control the number of threads assigned to each one, including the default queue:

- `--default-parallelism=N` specifies the default number of threads used for each queue. In the absence of this option, *N* is 2.

  The default queue always uses the default number of threads. Additional queues use the default number of threads unless you specify otherwise.

- `--parallel-schemas=[N:]db_list` sets up a processing queue for dumping the databases named in *`db_list`* and optionally specifies how many threads the queue uses. *`db_list`* is a comma-separated list of database names. If the option argument begins with *`N`*`:`, the queue uses *`N`* threads. Otherwise, the `--default-parallelism` option determines the number of queue threads.

  Multiple instances of the `--parallel-schemas` option create multiple queues.

  Names in the database list are permitted to contain the same `%` and `_` wildcard characters supported for filtering options (see mysqlpump Object Selection).

`mysqlpump` uses the default queue for processing any databases not named explicitly with a `--parallel-schemas` option, and for dumping user definitions if command options select them.

In general, with multiple queues, `mysqlpump` uses parallelism between the sets of databases processed by the queues, to dump multiple databases simultaneously. For a queue that uses multiple threads, `mysqlpump` uses parallelism within databases, to dump multiple objects from a given database simultaneously. Exceptions can occur; for example, `mysqlpump` may block queues while it obtains from the server lists of objects in databases.

With parallelism enabled, it is possible for output from different databases to be interleaved. For example, `INSERT` statements from multiple tables dumped in parallel can be interleaved; the statements are not written in any particular order. This does not affect reloading because output statements qualify object names with database names or are preceded by `USE` statements as required.

The granularity for parallelism is a single database object. For example, a single table cannot be dumped in parallel using multiple threads.

Examples:

```
shell> mysqlpump --parallel-schemas=db1,db2 --parallel-schemas=db3
```

`mysqlpump` sets up a queue to process `db1` and `db2`, another queue to process `db3`, and a default queue to process all other databases. All queues use two threads.

```
shell> mysqlpump --parallel-schemas=db1,db2 --parallel-schemas=db3
         --default-parallelism=4
```

This is the same as the previous example except that all queues use four threads.

```
shell> mysqlpump --parallel-schemas=5:db1,db2 --parallel-schemas=3:db3
```

The queue for `db1` and `db2` uses five threads, the queue for `db3` uses three threads, and the default queue uses the default of two threads.

As a special case, with `--default-parallelism=0` and no `--parallel-schemas` options, `mysqlpump` runs as a single-threaded process and creates no queues.

mysqlpump Restrictions

`mysqlpump` does not dump the `performance_schema`, `ndbinfo`, or `sys` schema by default. To dump any of these, name them explicitly on the command line. You can also name them with the `--databases` or `--include-databases` option.

`mysqlpump` does not dump the `INFORMATION_SCHEMA` schema.

`mysqlpump` dumps user accounts in logical form using `CREATE USER` and `GRANT` statements (for example, when you use the `--include-users` or `--users` option). For this reason, dumps of the `mysql` system database do not by default include the grant tables that contain user definitions: `user`, `db`, `tables_priv`, `columns_priv`, `procs_priv`, or `proxies_priv`. To dump any of the grant tables, name the `mysql` database followed by the table names:

```
shell> mysqlpump mysql user db ...
```

4.5.7 mysqlshow — Display Database, Table, and Column Information

The mysqlshow client can be used to quickly see which databases exist, their tables, or a table's columns or indexes.

mysqlshow provides a command-line interface to several SQL SHOW statements. See Section 13.7.5, "SHOW Syntax". The same information can be obtained by using those statements directly. For example, you can issue them from the mysql client program.

Invoke mysqlshow like this:

```
shell> mysqlshow [options] [db_name [tbl_name [col_name]]]
```

- If no database is given, a list of database names is shown.
- If no table is given, all matching tables in the database are shown.
- If no column is given, all matching columns and column types in the table are shown.

The output displays only the names of those databases, tables, or columns for which you have some privileges.

If the last argument contains shell or SQL wildcard characters (*, ?, %, or _), only those names that are matched by the wildcard are shown. If a database name contains any underscores, those should be escaped with a backslash (some Unix shells require two) to get a list of the proper tables or columns. * and ? characters are converted into SQL % and _ wildcard characters. This might cause some confusion when you try to display the columns for a table with a _ in the name, because in this case, mysqlshow shows you only the table names that match the pattern. This is easily fixed by adding an extra % last on the command line as a separate argument.

mysqlshow supports the following options, which can be specified on the command line or in the [mysqlshow] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, "Using Option Files".

Table 4.14 mysqlshow Options

Format	Description
--bind-address	Use specified network interface to connect to MySQL Server
--compress	Compress all information sent between client and server
--count	Show the number of rows per table
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--default-character-set	Specify default character set
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--enable-cleartext-plugin	Enable cleartext authentication plugin
--help	Display help message and exit
--host	Connect to MySQL server on given host

Format	Description
--keys	Show table indexes
--login-path	Read login path options from .mylogin.cnf
--no-defaults	Read no option files
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--protocol	Connection protocol to use
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--show-table-type	Show a column indicating the table type
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--status	Display extra information about each table
--tls-version	Protocols permitted for secure connections
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version	Display version information and exit

- `--help`, `-?`

  Display a help message and exit.

- `--bind-address=`*`ip_address`*

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--count`

  Show the number of rows per table. This can be slow for non-`MyISAM` tables.

- `--debug[=debug_options]`, `-# [debug_options]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:o`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-character-set=charset_name`

  Use *`charset_name`* as the default character set. See Section 10.5, “Character Set Configuration”.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqlshow` normally reads the `[client]` and `[mysqlshow]` groups. If the `--defaults-group-suffix=_other` option is given, `mysqlshow` also reads the `[client_other]` and `[mysqlshow_other]` groups.

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--host=host_name`, `-h host_name`

  Connect to the MySQL server on the given host.

- `--keys`, `-k`

  Show table indexes.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--no-defaults`

Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, "`mysql_config_editor` — MySQL Configuration Utility".)

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysqlshow` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, "End-User Guidelines for Password Security". You can use an option file to avoid giving the password on the command line.

- `--pipe`, `-W`

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqlshow` does not find it. See Section 6.3.10, "Pluggable Authentication".

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for the connection.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, "Connecting to the MySQL Server".

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--shared-memory-base-name=name`

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--show-table-type`, `-t`

  Show a column indicating the table type, as in `SHOW FULL TABLES`. The type is `BASE TABLE` or `VIEW`.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--status`, `-i`

  Display extra information about each table.

- `--tls-version=protocol_list`

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does. This option can be used multiple times to increase the amount of information.

- `--version`, `-V`

  Display version information and exit.

4.5.8 `mysqlslap` — Load Emulation Client

`mysqlslap` is a diagnostic program designed to emulate client load for a MySQL server and to report the timing of each stage. It works as if multiple clients are accessing the server.

Invoke `mysqlslap` like this:

```
shell> mysqlslap [options]
```

Some options such as `--create` or `--query` enable you to specify a string containing an SQL statement or a file containing statements. If you specify a file, by default it must contain one statement per line. (That is, the implicit statement delimiter is the newline character.) Use the `--delimiter` option to specify a different delimiter, which enables you to specify statements that span multiple lines or place multiple statements on a single line. You cannot include comments in a file; `mysqlslap` does not understand them.

`mysqlslap` runs in three stages:

1. Create schema, table, and optionally any stored programs or data to use for the test. This stage uses a single client connection.
2. Run the load test. This stage can use many client connections.
3. Clean up (disconnect, drop table if specified). This stage uses a single client connection.

Examples:

Supply your own create and query SQL statements, with 50 clients querying and 200 selects for each (enter the command on a single line):

```
mysqlslap --delimiter=";"
  --create="CREATE TABLE a (b int);INSERT INTO a VALUES (23)"
  --query="SELECT * FROM a" --concurrency=50 --iterations=200
```

Let `mysqlslap` build the query SQL statement with a table of two `INT` columns and three `VARCHAR` columns. Use five clients querying 20 times each. Do not create the table or insert the data (that is, use the previous test's schema and data):

```
mysqlslap --concurrency=5 --iterations=20
  --number-int-cols=2 --number-char-cols=3
  --auto-generate-sql
```

Tell the program to load the create, insert, and query SQL statements from the specified files, where the `create.sql` file has multiple table creation statements delimited by `';'` and multiple insert statements delimited by `';'`. The `--query` file will have multiple queries delimited by `';'`. Run all the load statements, then run all the queries in the query file with five clients (five times each):

```
mysqlslap --concurrency=5
  --iterations=5 --query=query.sql --create=create.sql
  --delimiter=";"
```

`mysqlslap` supports the following options, which can be specified on the command line or in the `[mysqlslap]` and `[client]` groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.15 `mysqlslap` Options

Format	Description
--auto-generate-sql	Generate SQL statements automatically when they are not supplied in files or using command options
--auto-generate-sql-add-autoincrement	Add AUTO_INCREMENT column to automatically generated tables
--auto-generate-sql-execute-number	Specify how many queries to generate automatically
--auto-generate-sql-guid-primary	Add a GUID-based primary key to automatically generated tables
--auto-generate-sql-load-type	Specify the test load type
--auto-generate-sql-secondary-indexes	Specify how many secondary indexes to add to automatically generated tables
--auto-generate-sql-unique-query-number	How many different queries to generate for automatic tests.
--auto-generate-sql-unique-write-number	How many different queries to generate for --auto-generate-sql-write-number
--auto-generate-sql-write-number	How many row inserts to perform on each thread
--commit	How many statements to execute before committing.
--compress	Compress all information sent between client and server

Format	Description
--concurrency	Number of clients to simulate when issuing the SELECT statement
--create	File or string containing the statement to use for creating the table
--create-schema	Schema in which to run the tests
--csv	Generate output in comma-separated values format
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--delimiter	Delimiter to use in SQL statements
--detach	Detach (close and reopen) each connection after each N statements
--enable-cleartext-plugin	Enable cleartext authentication plugin
--engine	Storage engine to use for creating the table
--help	Display help message and exit
--host	Connect to MySQL server on given host
--iterations	Number of times to run the tests
--login-path	Read login path options from .mylogin.cnf
--no-defaults	Read no option files
--no-drop	Do not drop any schema created during the test run
--number-char-cols	Number of VARCHAR columns to use if --auto-generate-sql is specified
--number-int-cols	Number of INT columns to use if --auto-generate-sql is specified
--number-of-queries	Limit each client to approximately this number of queries
--only-print	Do not connect to databases. mysqlslap only prints what it would have done
--password	Password to use when connecting to server
--pipe	On Windows, connect to server using named pipe
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--post-query	File or string containing the statement to execute after the tests have completed
--post-system	String to execute using system() after the tests have completed
--pre-query	File or string containing the statement to execute before running the tests
--pre-system	String to execute using system() before running the tests
--print-defaults	Print default options
--protocol	Connection protocol to use

Format	Description
--query	File or string containing the SELECT statement to use for retrieving data
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--silent	Silent mode
--socket	For connections to localhost, the Unix socket file to use
--sql-mode	Set SQL mode for client session
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tls-version	Protocols permitted for secure connections
--user	MySQL user name to use when connecting to server
--verbose	Verbose mode
--version	Display version information and exit

- `--help`, `-?`

  Display a help message and exit.

- `--auto-generate-sql`, `-a`

  Generate SQL statements automatically when they are not supplied in files or using command options.

- `--auto-generate-sql-add-autoincrement`

  Add an `AUTO_INCREMENT` column to automatically generated tables.

- `--auto-generate-sql-execute-number=`*`N`*

  Specify how many queries to generate automatically.

- `--auto-generate-sql-guid-primary`

  Add a GUID-based primary key to automatically generated tables.

- `--auto-generate-sql-load-type=`*`type`*

  Specify the test load type. The permissible values are `read` (scan tables), `write` (insert into tables), `key` (read primary keys), `update` (update primary keys), or `mixed` (half inserts, half scanning selects). The default is `mixed`.

- `--auto-generate-sql-secondary-indexes=`*`N`*

  Specify how many secondary indexes to add to automatically generated tables. By default, none are added.

- `--auto-generate-sql-unique-query-number=N`

  How many different queries to generate for automatic tests. For example, if you run a `key` test that performs 1000 selects, you can use this option with a value of 1000 to run 1000 unique queries, or with a value of 50 to perform 50 different selects. The default is 10.

- `--auto-generate-sql-unique-write-number=N`

  How many different queries to generate for `--auto-generate-sql-write-number`. The default is 10.

- `--auto-generate-sql-write-number=N`

  How many row inserts to perform. The default is 100.

- `--commit=N`

  How many statements to execute before committing. The default is 0 (no commits are done).

- `--compress`, `-C`

  Compress all information sent between the client and the server if both support compression.

- `--concurrency=N`, `-c N`

  The number of parallel clients to simulate.

- `--create=value`

  The file or string containing the statement to use for creating the table.

- `--create-schema=value`

  The schema in which to run the tests.

  Note

  If the `--auto-generate-sql` option is also given, `mysqlslap` drops the schema at the end of the test run. To avoid this, use the `--no-drop` option as well.

- `--csv[=file_name]`

  Generate output in comma-separated values format. The output goes to the named file, or to the standard output if no file is given.

- `--debug[=debug_options]`, `-# [debug_options]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:o,/tmp/mysqlslap.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`, `-T`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

  Exception: Even with `--defaults-file`, client programs read `.mylogin.cnf`.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqlslap` normally reads the `[client]` and `[mysqlslap]` groups. If the `--defaults-group-suffix=_other` option is given, `mysqlslap` also reads the `[client_other]` and `[mysqlslap_other]` groups.

- `--delimiter=str`, `-F str`

  The delimiter to use in SQL statements supplied in files or using command options.

- `--detach=N`

  Detach (close and reopen) each connection after each *`N`* statements. The default is 0 (connections are not detached).

- `--enable-cleartext-plugin`

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `--engine=engine_name`, `-e engine_name`

  The storage engine to use for creating tables.

- `--host=host_name`, `-h host_name`

  Connect to the MySQL server on the given host.

- `--iterations=N`, `-i N`

  The number of times to run the tests.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--no-drop`

  Prevent `mysqlslap` from dropping any schema it creates during the test run.

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

The exception is that the .mylogin.cnf file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when --no-defaults is used. (.mylogin.cnf is created by the mysql_config_editor utility. See Section 4.6.7, "mysql_config_editor — MySQL Configuration Utility".)

- --number-char-cols=*N*, -x *N*

  The number of VARCHAR columns to use if --auto-generate-sql is specified.

- --number-int-cols=*N*, -y *N*

  The number of INT columns to use if --auto-generate-sql is specified.

- --number-of-queries=*N*

  Limit each client to approximately this many queries. Query counting takes into account the statement delimiter. For example, if you invoke mysqlslap as follows, the ; delimiter is recognized so that each instance of the query string counts as two queries. As a result, 5 rows (not 10) are inserted.

```
shell> mysqlslap --delimiter=";" --number-of-queries=10
         --query="use test;insert into t values(null)"
```

- --only-print

  Do not connect to databases. mysqlslap only prints what it would have done.

- --password[=*password*], -p[*password*]

  The password to use when connecting to the server. If you use the short option form (-p), you *cannot* have a space between the option and the password. If you omit the *password* value following the --password or -p option on the command line, mysqlslap prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, "End-User Guidelines for Password Security". You can use an option file to avoid giving the password on the command line.

- --pipe, -W

  On Windows, connect to the server using a named pipe. This option applies only if the server supports named-pipe connections.

- --plugin-dir=*dir_name*

  The directory in which to look for plugins. Specify this option if the --default-auth option is used to specify an authentication plugin but mysqlslap does not find it. See Section 6.3.10, "Pluggable Authentication".

- --port=*port_num*, -P *port_num*

  The TCP/IP port number to use for the connection.

- --post-query=*value*

  The file or string containing the statement to execute after the tests have completed. This execution is not counted for timing purposes.

- --post-system=*str*

  The string to execute using system() after the tests have completed. This execution is not counted for timing purposes.

- `--pre-query=value`

  The file or string containing the statement to execute before running the tests. This execution is not counted for timing purposes.

- `--pre-system=str`

  The string to execute using `system()` before running the tests. This execution is not counted for timing purposes.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, “Connecting to the MySQL Server”.

- `--query=value`, `-q value`

  The file or string containing the `SELECT` statement to use for retrieving data.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--shared-memory-base-name=name`

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. This option applies only if the server supports shared-memory connections.

- `--silent`, `-s`

  Silent mode. No output.

- `--socket=path`, `-S path`

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--sql-mode=mode`

  Set the SQL mode for the client session.

- `--ssl*`

  Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--tls-version=protocol_list`

The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to the server.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does. This option can be used multiple times to increase the amount of information.

- `--version`, `-V`

  Display version information and exit.

4.6 MySQL Administrative and Utility Programs

This section describes administrative programs and programs that perform miscellaneous utility operations.

4.6.1 `ibd2sdi` — InnoDB Tablespace SDI Extraction Utility

`ibd2sdi` is a utility for extracting serialized dictionary information (SDI) from `InnoDB` tablespace files. SDI data is present in persistent `InnoDB` tablespace files as of MySQL 8.0.2.

`ibd2sdi` can be run on file-per-table tablespace files (`*.ibd` files), general tablespace files (`*.ibd` files), and system tablespace files (`ibdata*` files). `ibd2sdi` is not supported for use with the global data dictionary tablespace, the shared temporary tablespace, or undo log tablespaces.

`ibd2sdi` can be used at runtime or while the server is offline. During DDL operations, `ROLLBACK`, or undo log purge operations related to SDI records, there may be a short interval of time when `ibd2sdi` fails to read SDI data stored in the tablespace.

`ibd2sdi` performs an uncommitted read of SDI records from the specified tabelespace. Redo logs and undo logs are not accessed.

Invoke the `ibd2sdi` utility like this:

```
shell> ibd2sdi [options] file_name1 [file_name2 file_name3 ...]
```

`ibd2sdi` supports multi-file tablespaces like the `InnoDB` system tablespace, but it cannot be run on more than one tablespace at a time. For multi-file tablespaces, specify each file:

```
shell> ibd2sdi ibdata1 ibdata2
```

The files of a multi-file tablespace must be specified in order of the ascending page number. If two successive files have the same space ID, the later file must start with the last page number of the previous file + 1.

`ibd2sdi` outputs SDI records (containing id, type, and data fields) in `JSON` format.

ibd2sdi Options

`ibd2sdi` supports the following options:

- `--help`, `-h`

Displays command line help. Example usage:

```
shell> ibd2sdi --help
```

- `--version`, `-v`

  Displays MySQL version information. Example usage:

  ```
  shell> ibd2sdi --version
  ```

- `--dump-file=`, `-d`

  Dumps serialized dictionary information (SDI) into the specified dump file. If a dump file is not specified, the tablespace SDI is dumped to `stdout`. Example usage:

  ```
  shell> ibd2sdi --dump-file=file_name ../data/test/t1.ibd
  ```

- `--read=#`, `-r #`

  Reads from the specified copy of the serialized dictionary information (SDI). Permissible values are `0` or `1`. Example usage:

  ```
  shell> ibd2sdi --read=1 ../data/test/t1.ibd
  ```

- `--skip-data`, `-s`

  Skips retrieval of `data` field values from the serialized dictionary information (SDI) records and only retrieves the `id` and `type` field values, which are primary keys for SDI records. Example usage:

  ```
  shell> ibd2sdi --skip-data ../data/test/t1.ibd
  ```

- `--id=#`, `-i #`

  Retrieves serialized dictionary information (SDI) records matching the specified id. All records are scanned.

  ```
  shell> ibd2sdi --id=701 ../data/test/t1.ibd
  ```

- `--type=#`, `-t #`

  Retrieves serialized dictionary information (SDI) records matching the specified tablespace object type. All records are scanned.

  ```
  shell> ibd2sdi --type=3 ../data/test/t1.ibd
  ```

- `--strict-check`, `-c`

  Specifies a strict checksum algorithm for validating the checksum of pages that are read. Options include `innodb`, `crc32`, and `none`.

  In this example, the strict version of the `innodb` checksum algorithm is specified:

  ```
  shell> ibd2sdi --strict-check=innodb ../data/test/t1.ibd
  ```

  In this example, the strict version of `crc32` checksum algorithm is specified:

  ```
  shell> ibd2sdi -c crc32 ../data/test/t1.ibd
  ```

If you do not specify the `--strict-check` option, validation is performed against non-strict `innodb`, `crc32` and `none` checksums.

- `--no-check`, `-n`

  Skips checksum validation for pages that are read. Example usage:

```
shell> ibd2sdi --no-check ../data/test/t1.ibd
```

4.6.2 innochecksum — Offline InnoDB File Checksum Utility

`innochecksum` prints checksums for `InnoDB` files. This tool reads an `InnoDB` tablespace file, calculates the checksum for each page, compares the calculated checksum to the stored checksum, and reports mismatches, which indicate damaged pages. It was originally developed to speed up verifying the integrity of tablespace files after power outages but can also be used after file copies. Because checksum mismatches cause `InnoDB` to deliberately shut down a running server, it may be preferable to use this tool rather than waiting for an in-production server to encounter the damaged pages.

`innochecksum` cannot be used on tablespace files that the server already has open. For such files, you should use `CHECK TABLE` to check tables within the tablespace. Attempting to run `innochecksum` on a tablespace that the server already has open will result in an "`Unable to lock file`" error.

If checksum mismatches are found, you would normally restore the tablespace from backup or start the server and attempt to use `mysqldump` to make a backup of the tables within the tablespace.

Invoke `innochecksum` like this:

```
shell> innochecksum [options] file_name
```

innochecksum Options

`innochecksum` supports the following options. For options that refer to page numbers, the numbers are zero-based.

- `--help`, `-?`

  Displays command line help. Example usage:

```
shell> innochecksum --help
```

- `--info`, `-I`

  Synonym for `--help`. Displays command line help. Example usage:

```
shell> innochecksum --info
```

- `--version`, `-V`

  Displays version information. Example usage:

```
shell> innochecksum --version
```

- `--verbose`, `-v`

  Verbose mode; prints a progress indicator to the log file every five seconds. In order for the progress indicator to be printed, the log file must be specified using the `--log option`. To turn on `verbose` mode, run:

```
shell> innochecksum --verbose
```

To turn off verbose mode, run:

```
shell> innochecksum --verbose=FALSE
```

The `--verbose` option and `--log` option can be specified at the same time. For example:

```
shell> innochecksum --verbose --log=/var/lib/mysql/test/logtest.txt
```

To locate the progress indicator information in the log file, you can perform the following search:

```
shell> cat ./logtest.txt | grep -i "okay"
```

The progress indicator information in the log file appears similar to the following:

```
page 1663 okay: 2.863% done
page 8447 okay: 14.537% done
page 13695 okay: 23.568% done
page 18815 okay: 32.379% done
page 23039 okay: 39.648% done
page 28351 okay: 48.789% done
page 33023 okay: 56.828% done
page 37951 okay: 65.308% done
page 44095 okay: 75.881% done
page 49407 okay: 85.022% done
page 54463 okay: 93.722% done
...
```

- `--count`, `-c`

  Print a count of the number of pages in the file and exit. Example usage:

  ```
  shell> innochecksum --count ../data/test/tab1.ibd
  ```

- `--start-page=num`, `-s num`

  Start at this page number. Example usage:

  ```
  shell> innochecksum --start-page=600 ../data/test/tab1.ibd
  ```

  or:

  ```
  shell> innochecksum -s 600 ../data/test/tab1.ibd
  ```

- `--end-page=num`, `-e num`

  End at this page number. Example usage:

  ```
  shell> innochecksum --end-page=700 ../data/test/tab1.ibd
  ```

  or:

  ```
  shell> innochecksum --p 700 ../data/test/tab1.ibd
  ```

- `--page=num`, `-p num`

  Check only this page number. Example usage:

```
shell> innochecksum --page=701 ../data/test/tab1.ibd
```

- `--strict-check`, `-C`

  Specify a strict checksum algorithm. Options include `innodb`, `crc32`, and `none`.

  In this example, the `innodb` checksum algorithm is specified:

  ```
  shell> innochecksum --strict-check=innodb ../data/test/tab1.ibd
  ```

  In this example, the `crc32` checksum algorithm is specified:

  ```
  shell> innochecksum -C crc32 ../data/test/tab1.ibd
  ```

  The following conditions apply:

  - If you do not specify the `--strict-check` option, `innochecksum` validates against `innodb`, `crc32` and `none`.
  - If you specify the `none` option, only checksums generated by `none` are allowed.
  - If you specify the `innodb` option, only checksums generated by `innodb` are allowed.
  - If you specify the `crc32` option, only checksums generated by `crc32` are allowed.

- `--no-check`, `-n`

  Ignore the checksum verification when rewriting a checksum. This option may only be used with the `innochecksum --write` option. If the `--write` option is not specified, `innochecksum` will terminate.

  In this example, an `innodb` checksum is rewritten to replace an invalid checksum:

  ```
  shell> innochecksum --no-check --write innodb ../data/test/tab1.ibd
  ```

- `--allow-mismatches`, `-a`

  The maximum number of checksum mismatches allowed before `innochecksum` terminates. The default setting is 0. If `--allow-mismatches=N`, where `N>=0`, `N` mismatches are permitted and `innochecksum` terminates at `N+1`. When `--allow-mismatches` is set to 0, `innochecksum` terminates on the first checksum mismatch.

  In this example, an existing `innodb` checksum is rewritten to set `--allow-mismatches` to 1.

  ```
  shell> innochecksum --allow-mismatches=1 --write innodb ../data/test/tab1.ibd
  ```

  With `--allow-mismatches` set to 1, if there is a mismatch at page 600 and another at page 700 on a file with 1000 pages, the checksum is updated for pages 0-599 and 601-699. Because `--allow-mismatches` is set to 1, the checksum tolerates the first mismatch and terminates on the second mismatch, leaving page 600 and pages 700-999 unchanged.

- `--write=name`, `-w num`

  Rewrite a checksum. When rewriting an invalid checksum, the `--no-check` option must be used together with the `--write` option. The `--no-check` option tells `innochecksum` to ignore verification of the invalid checksum. You do not have to specify the `--no-check` option if the current checksum is valid.

An algorithm must be specified when using the `--write` option. Possible values for the `--write` option are:

- `innodb`: A checksum calculated in software, using the original algorithm from `InnoDB`.
- `crc32`: A checksum calculated using the `crc32` algorithm, possibly done with a hardware assist.
- `none`: A constant number.

The `--write` option rewrites entire pages to disk. If the new checksum is identical to the existing checksum, the new checksum is not written to disk in order to minimize I/O.

`innochecksum` obtains an exclusive lock when the `--write` option is used.

In this example, a `crc32` checksum is written for `tab1.ibd`:

```
shell> innochecksum -w crc32 ../data/test/tab1.ibd
```

In this example, a `crc32` checksum is rewritten to replace an invalid `crc32` checksum:

```
shell> innochecksum --no-check --write crc32 ../data/test/tab1.ibd
```

- `--page-type-summary`, `-S`

  Display a count of each page type in a tablespace. Example usage:

  ```
  shell> innochecksum --page-type-summary ../data/test/tab1.ibd
  ```

  Sample output for `--page-type-summary`:

  ```
  File::../data/test/tab1.ibd
  ================PAGE TYPE SUMMARY==============
  #PAGE_COUNT PAGE_TYPE
  ===============================================
        2        Index page
        0        Undo log page
        1        Inode page
        0        Insert buffer free list page
        2        Freshly allocated page
        1        Insert buffer bitmap
        0        System page
        0        Transaction system page
        1        File Space Header
        0        Extent descriptor page
        0        BLOB page
        0        Compressed BLOB page
        0        Other type of page
  ===============================================
  Additional information:
  Undo page type: 0 insert, 0 update, 0 other
  Undo page state: 0 active, 0 cached, 0 to_free, 0 to_purge, 0 prepared, 0 other
  ```

- `--page-type-dump`, `-D`

  Dump the page type information for each page in a tablespace to `stderr` or `stdout`. Example usage:

  ```
  shell> innochecksum --page-type-dump=/tmp/a.txt ../data/test/tab1.ibd
  ```

- `--log`, `-l`

Log output for the innochecksum tool. A log file name must be provided. Log output contains checksum values for each tablespace page. For uncompressed tables, LSN values are also provided. The --log replaces the --debug option, which was available in earlier releases. Example usage:

```
shell> innochecksum --log=/tmp/log.txt ../data/test/tab1.ibd
```

or:

```
shell> innochecksum -l /tmp/log.txt ../data/test/tab1.ibd
```

- - option.

  Specify the - option to read from standard input. If the - option is missing when "read from standard in" is expected, innochecksum will output innochecksum usage information indicating that the "-" option was omitted. Example usages:

  ```
  shell> cat t1.ibd | innochecksum -
  ```

  In this example, innochecksum writes the crc32 checksum algorithm to a.ibd without changing the original t1.ibd file.

  ```
  shell> cat t1.ibd | innochecksum --write=crc32 - > a.ibd
  ```

Running innochecksum on Multiple User-defined Tablespace Files

The following examples demonstrate how to run innochecksum on multiple user-defined tablespace files (.ibd files).

Run innochecksum for all tablespace (.ibd) files in the "test" database:

```
shell> innochecksum ./data/test/*.ibd
```

Run innochecksum for all tablespace files (.ibd files) that have a file name starting with "t":

```
shell> innochecksum ./data/test/t*.ibd
```

Run innochecksum for all tablespace files (.ibd files) in the data directory:

```
shell> innochecksum ./data/*/*.ibd
```

Note

Running innochecksum on multiple user-defined tablespace files is not supported on Windows operating systems, as Windows shells such as cmd.exe do not support glob pattern expansion. On Windows systems, innochecksum must be run separately for each user-defined tablespace file. For example:

```
cmd> innochecksum.exe t1.ibd
cmd> innochecksum.exe t2.ibd
cmd> innochecksum.exe t3.ibd
```

Running innochecksum on Multiple System Tablespace Files

By default, there is only one InnoDB system tablespace file (ibdata1) but multiple files for the system tablespace can be defined using the innodb_data_file_path option. In the following example, three files for the system tablespace are defined using the innodb_data_file_path option: ibdata1, ibdata2, and ibdata3.

```
shell> ./bin/mysqld --no-defaults --innodb-data-file-path="ibdata1:10M;ibdata2:10M;ibdata3:10M:autoexte
```

The three files (ibdata1, ibdata2, and ibdata3) form one logical system tablespace. To run innochecksum on multiple files that form one logical system tablespace, innochecksum requires the - option to read tablespace files in from standard input, which is equivalent to concatenating multiple files to create one single file. For the example provided above, the following innochecksum command would be used:

```
shell> cat ibdata* | innochecksum -
```

Refer to the innochecksum options information for more information about the "-" option.

Note

Running innochecksum on multiple files in the same tablespace is not supported on Windows operating systems, as Windows shells such as cmd.exe do not support glob pattern expansion. On Windows systems, innochecksum must be run separately for each system tablespace file. For example:

```
cmd> innochecksum.exe ibdata1
cmd> innochecksum.exe ibdata2
cmd> innochecksum.exe ibdata3
```

4.6.3 myisam_ftdump — Display Full-Text Index information

myisam_ftdump displays information about FULLTEXT indexes in MyISAM tables. It reads the MyISAM index file directly, so it must be run on the server host where the table is located. Before using myisam_ftdump, be sure to issue a FLUSH TABLES statement first if the server is running.

myisam_ftdump scans and dumps the entire index, which is not particularly fast. On the other hand, the distribution of words changes infrequently, so it need not be run often.

Invoke myisam_ftdump like this:

```
shell> myisam_ftdump [options] tbl_name index_num
```

The *tbl_name* argument should be the name of a MyISAM table. You can also specify a table by naming its index file (the file with the .MYI suffix). If you do not invoke myisam_ftdump in the directory where the table files are located, the table or index file name must be preceded by the path name to the table's database directory. Index numbers begin with 0.

Example: Suppose that the test database contains a table named mytexttable that has the following definition:

```
CREATE TABLE mytexttable
(
  id   INT NOT NULL,
  txt  TEXT NOT NULL,
  PRIMARY KEY (id),
  FULLTEXT (txt)
) ENGINE=MyISAM;
```

The index on id is index 0 and the FULLTEXT index on txt is index 1. If your working directory is the test database directory, invoke myisam_ftdump as follows:

```
shell> myisam_ftdump mytexttable 1
```

If the path name to the `test` database directory is `/usr/local/mysql/data/test`, you can also specify the table name argument using that path name. This is useful if you do not invoke `myisam_ftdump` in the database directory:

```
shell> myisam_ftdump /usr/local/mysql/data/test/mytexttable 1
```

You can use `myisam_ftdump` to generate a list of index entries in order of frequency of occurrence like this:

```
shell> myisam_ftdump -c mytexttable 1 | sort -r
```

`myisam_ftdump` supports the following options:

- `--help`, `-h -?`

  Display a help message and exit.

- `--count`, `-c`

  Calculate per-word statistics (counts and global weights).

- `--dump`, `-d`

  Dump the index, including data offsets and word weights.

- `--length`, `-l`

  Report the length distribution.

- `--stats`, `-s`

  Report global index statistics. This is the default operation if no other operation is specified.

- `--verbose`, `-v`

  Verbose mode. Print more output about what the program does.

4.6.4 `myisamchk` — MyISAM Table-Maintenance Utility

The `myisamchk` utility gets information about your database tables or checks, repairs, or optimizes them. `myisamchk` works with `MyISAM` tables (tables that have `.MYD` and `.MYI` files for storing data and indexes).

You can also use the `CHECK TABLE` and `REPAIR TABLE` statements to check and repair `MyISAM` tables. See Section 13.7.2.2, “CHECK TABLE Syntax”, and Section 13.7.2.5, “REPAIR TABLE Syntax”.

The use of `myisamchk` with partitioned tables is not supported.

Caution

It is best to make a backup of a table before performing a table repair operation; under some circumstances the operation might cause data loss. Possible causes include but are not limited to file system errors.

Invoke `myisamchk` like this:

```
shell> myisamchk [options] tbl_name ...
```

The *`options`* specify what you want `myisamchk` to do. They are described in the following sections. You can also get a list of options by invoking `myisamchk --help`.

With no options, myisamchk simply checks your table as the default operation. To get more information or to tell myisamchk to take corrective action, specify options as described in the following discussion.

tbl_name is the database table you want to check or repair. If you run myisamchk somewhere other than in the database directory, you must specify the path to the database directory, because myisamchk has no idea where the database is located. In fact, myisamchk does not actually care whether the files you are working on are located in a database directory. You can copy the files that correspond to a database table into some other location and perform recovery operations on them there.

You can name several tables on the myisamchk command line if you wish. You can also specify a table by naming its index file (the file with the .MYI suffix). This enables you to specify all tables in a directory by using the pattern *.MYI. For example, if you are in a database directory, you can check all the MyISAM tables in that directory like this:

```
shell> myisamchk *.MYI
```

If you are not in the database directory, you can check all the tables there by specifying the path to the directory:

```
shell> myisamchk /path/to/database_dir/*.MYI
```

You can even check all tables in all databases by specifying a wildcard with the path to the MySQL data directory:

```
shell> myisamchk /path/to/datadir/*/*.MYI
```

The recommended way to quickly check all MyISAM tables is:

```
shell> myisamchk --silent --fast /path/to/datadir/*/*.MYI
```

If you want to check all MyISAM tables and repair any that are corrupted, you can use the following command:

```
shell> myisamchk --silent --force --fast --update-state \
          --key_buffer_size=64M --myisam_sort_buffer_size=64M \
          --read_buffer_size=1M --write_buffer_size=1M \
          /path/to/datadir/*/*.MYI
```

This command assumes that you have more than 64MB free. For more information about memory allocation with myisamchk, see Section 4.6.4.6, “myisamchk Memory Usage”.

For additional information about using myisamchk, see Section 7.6, “MyISAM Table Maintenance and Crash Recovery”.

Important

You must ensure that no other program is using the tables while you are running myisamchk. The most effective means of doing so is to shut down the MySQL server while running myisamchk, or to lock all tables that myisamchk is being used on.

Otherwise, when you run myisamchk, it may display the following error message:

```
warning: clients are using or haven't closed the table properly
```

This means that you are trying to check a table that has been updated by another program (such as the mysqld server) that hasn't yet closed the file or

> that has died without closing the file properly, which can sometimes lead to the corruption of one or more MyISAM tables.
>
> If mysqld is running, you must force it to flush any table modifications that are still buffered in memory by using FLUSH TABLES. You should then ensure that no one is using the tables while you are running myisamchk
>
> However, the easiest way to avoid this problem is to use CHECK TABLE instead of myisamchk to check tables. See Section 13.7.2.2, “CHECK TABLE Syntax”.

myisamchk supports the following options, which can be specified on the command line or in the [myisamchk] group of an option file. For information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Table 4.16 myisamchk Options

Format	Description
--analyze	Analyze the distribution of key values
--backup	Make a backup of the .MYD file as file_name-time.BAK
--block-search	Find the record that a block at the given offset belongs to
--check	Check the table for errors
--check-only-changed	Check only tables that have changed since the last check
--correct-checksum	Correct the checksum information for the table
--data-file-length	Maximum length of the data file (when re-creating data file when it is full)
--debug	Write debugging log
--decode_bits	Decode_bits
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--description	Print some descriptive information about the table
--extend-check	Do very thorough table check or repair that tries to recover every possible row from the data file
--fast	Check only tables that haven't been closed properly
--force	Do a repair operation automatically if myisamchk finds any errors in the table
--force	Overwrite old temporary files. For use with the -r or -o option
--ft_max_word_len	Maximum word length for FULLTEXT indexes
--ft_min_word_len	Minimum word length for FULLTEXT indexes
--ft_stopword_file	Use stopwords from this file instead of built-in list
--HELP	Display help message and exit
--help	Display help message and exit
--information	Print informational statistics about the table that is checked
--key_buffer_size	Size of buffer used for index blocks for MyISAM tables
--keys-used	A bit-value that indicates which indexes to update
--max-record-length	Skip rows larger than the given length if myisamchk cannot allocate memory to hold them
--medium-check	Do a check that is faster than an --extend-check operation

Format	Description
--myisam_block_size	Block size to be used for MyISAM index pages
--myisam_sort_buffer_size	The buffer that is allocated when sorting the index when doing a REPAIR or when creating indexes with CREATE INDEX or ALTER TABLE
--no-defaults	Read no option files
--parallel-recover	Uses the same technique as -r and -n, but creates all the keys in parallel, using different threads (beta)
--print-defaults	Print default options
--quick	Achieve a faster repair by not modifying the data file.
--read_buffer_size	Each thread that does a sequential scan allocates a buffer of this size for each table it scans
--read-only	Do not mark the table as checked
--recover	Do a repair that can fix almost any problem except unique keys that aren't unique
--safe-recover	Do a repair using an old recovery method that reads through all rows in order and updates all index trees based on the rows found
--set-auto-increment	Force AUTO_INCREMENT numbering for new records to start at the given value
--set-collation	Specify the collation to use for sorting table indexes
--silent	Silent mode
--sort_buffer_size	The buffer that is allocated when sorting the index when doing a REPAIR or when creating indexes with CREATE INDEX or ALTER TABLE
--sort-index	Sort the index tree blocks in high-low order
--sort_key_blocks	sort_key_blocks
--sort-records	Sort records according to a particular index
--sort-recover	Force myisamchk to use sorting to resolve the keys even if the temporary files would be very large
--stats_method	Specifies how MyISAM index statistics collection code should treat NULLs
--tmpdir	Path of the directory to be used for storing temporary files
--unpack	Unpack a table that was packed with myisampack
--update-state	Store information in the .MYI file to indicate when the table was checked and whether the table crashed
--verbose	Verbose mode
--version	Display version information and exit
--write_buffer_size	Write buffer size

4.6.4.1 myisamchk General Options

The options described in this section can be used for any type of table maintenance operation performed by `myisamchk`. The sections following this one describe options that pertain only to specific operations, such as table checking or repairing.

- `--help`, `-?`

  Display a help message and exit. Options are grouped by type of operation.

- `--HELP`, `-H`

  Display a help message and exit. Options are presented in a single list.

- `--debug=debug_options`, `-# debug_options`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:o,/tmp/myisamchk.trace`.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

  Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-group-suffix=str`

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `myisamchk` normally reads the `[myisamchk]` group. If the `--defaults-group-suffix=_other` option is given, `myisamchk` also reads the `[myisamchk_other]` group.

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.)

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--silent`, `-s`

  Silent mode. Write output only when errors occur. You can use `-s` twice (`-ss`) to make `myisamchk` very silent.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does. This can be used with `-d` and `-e`. Use `-v` multiple times (`-vv`, `-vvv`) for even more output.

- `--version`, `-V`

  Display version information and exit.

- `--wait`, `-w`

  Instead of terminating with an error if the table is locked, wait until the table is unlocked before continuing. If you are running `mysqld` with external locking disabled, the table can be locked only by another `myisamchk` command.

You can also set the following variables by using `--var_name=value` syntax:

Variable	Default Value
decode_bits	9
ft_max_word_len	version-dependent
ft_min_word_len	4
ft_stopword_file	built-in list
key_buffer_size	523264
myisam_block_size	1024
myisam_sort_key_blocks	16
read_buffer_size	262136
sort_buffer_size	2097144
sort_key_blocks	16
stats_method	nulls_unequal
write_buffer_size	262136

The possible myisamchk variables and their default values can be examined with myisamchk --help:

myisam_sort_buffer_size is used when the keys are repaired by sorting keys, which is the normal case when you use --recover. sort_buffer_size is a deprecated synonym for myisam_sort_buffer_size.

key_buffer_size is used when you are checking the table with --extend-check or when the keys are repaired by inserting keys row by row into the table (like when doing normal inserts). Repairing through the key buffer is used in the following cases:

- You use --safe-recover.
- The temporary files needed to sort the keys would be more than twice as big as when creating the key file directly. This is often the case when you have large key values for CHAR, VARCHAR, or TEXT columns, because the sort operation needs to store the complete key values as it proceeds. If you have lots of temporary space and you can force myisamchk to repair by sorting, you can use the --sort-recover option.

Repairing through the key buffer takes much less disk space than using sorting, but is also much slower.

If you want a faster repair, set the key_buffer_size and myisam_sort_buffer_size variables to about 25% of your available memory. You can set both variables to large values, because only one of them is used at a time.

myisam_block_size is the size used for index blocks.

stats_method influences how NULL values are treated for index statistics collection when the --analyze option is given. It acts like the myisam_stats_method system variable. For more information, see the description of myisam_stats_method in Section 5.1.5, “Server System Variables”, and Section 8.3.7, “InnoDB and MyISAM Index Statistics Collection”.

ft_min_word_len and ft_max_word_len indicate the minimum and maximum word length for FULLTEXT indexes on MyISAM tables. ft_stopword_file names the stopword file. These need to be set under the following circumstances.

If you use myisamchk to perform an operation that modifies table indexes (such as repair or analyze), the FULLTEXT indexes are rebuilt using the default full-text parameter values for minimum and maximum word length and the stopword file unless you specify otherwise. This can result in queries failing.

The problem occurs because these parameters are known only by the server. They are not stored in MyISAM index files. To avoid the problem if you have modified the minimum or maximum word length or the stopword file in the server, specify the same ft_min_word_len, ft_max_word_len, and ft_stopword_file values to myisamchk that you use for mysqld. For example, if you have set the minimum word length to 3, you can repair a table with myisamchk like this:

```
shell> myisamchk --recover --ft_min_word_len=3 tbl_name.MYI
```

To ensure that myisamchk and the server use the same values for full-text parameters, you can place each one in both the [mysqld] and [myisamchk] sections of an option file:

```
[mysqld]
ft_min_word_len=3

[myisamchk]
ft_min_word_len=3
```

An alternative to using myisamchk is to use the REPAIR TABLE, ANALYZE TABLE, OPTIMIZE TABLE, or ALTER TABLE. These statements are performed by the server, which knows the proper full-text parameter values to use.

4.6.4.2 myisamchk Check Options

myisamchk supports the following options for table checking operations:

- --check, -c

  Check the table for errors. This is the default operation if you specify no option that selects an operation type explicitly.

- --check-only-changed, -C

  Check only tables that have changed since the last check.

- --extend-check, -e

  Check the table very thoroughly. This is quite slow if the table has many indexes. This option should only be used in extreme cases. Normally, myisamchk or myisamchk --medium-check should be able to determine whether there are any errors in the table.

  If you are using --extend-check and have plenty of memory, setting the key_buffer_size variable to a large value helps the repair operation run faster.

  See also the description of this option under table repair options.

  For a description of the output format, see Section 4.6.4.5, “Obtaining Table Information with myisamchk”.

- --fast, -F

  Check only tables that haven't been closed properly.

- --force, -f

  Do a repair operation automatically if myisamchk finds any errors in the table. The repair type is the same as that specified with the --recover or -r option.

- --information, -i

  Print informational statistics about the table that is checked.

- --medium-check, -m

Do a check that is faster than an `--extend-check` operation. This finds only 99.99% of all errors, which should be good enough in most cases.

- `--read-only`, `-T`

  Do not mark the table as checked. This is useful if you use `myisamchk` to check a table that is in use by some other application that does not use locking, such as `mysqld` when run with external locking disabled.

- `--update-state`, `-U`

  Store information in the `.MYI` file to indicate when the table was checked and whether the table crashed. This should be used to get full benefit of the `--check-only-changed` option, but you shouldn't use this option if the `mysqld` server is using the table and you are running it with external locking disabled.

4.6.4.3 myisamchk Repair Options

`myisamchk` supports the following options for table repair operations (operations performed when an option such as `--recover` or `--safe-recover` is given):

- `--backup`, `-B`

  Make a backup of the `.MYD` file as *`file_name-time`*`.BAK`

- `--character-sets-dir=`*`dir_name`*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--correct-checksum`

  Correct the checksum information for the table.

- `--data-file-length=`*`len`*, `-D` *`len`*

  The maximum length of the data file (when re-creating data file when it is “full”).

- `--extend-check`, `-e`

  Do a repair that tries to recover every possible row from the data file. Normally, this also finds a lot of garbage rows. Do not use this option unless you are desperate.

  See also the description of this option under table checking options.

  For a description of the output format, see Section 4.6.4.5, “Obtaining Table Information with myisamchk”.

- `--force`, `-f`

  Overwrite old intermediate files (files with names like *`tbl_name`*`.TMD`) instead of aborting.

- `--keys-used=`*`val`*, `-k` *`val`*

  For `myisamchk`, the option value is a bit value that indicates which indexes to update. Each binary bit of the option value corresponds to a table index, where the first index is bit 0. An option value of 0 disables updates to all indexes, which can be used to get faster inserts. Deactivated indexes can be reactivated by using `myisamchk -r`.

- `--no-symlinks`, `-l`

  Do not follow symbolic links. Normally `myisamchk` repairs the table that a symlink points to. This option does not exist as of MySQL 4.0 because versions from 4.0 on do not remove symlinks during repair operations.

- `--max-record-length=len`

  Skip rows larger than the given length if `myisamchk` cannot allocate memory to hold them.

- `--parallel-recover`, `-p`

  Use the same technique as `-r` and `-n`, but create all the keys in parallel, using different threads. *This is beta-quality code. Use at your own risk!*

- `--quick`, `-q`

  Achieve a faster repair by modifying only the index file, not the data file. You can specify this option twice to force `myisamchk` to modify the original data file in case of duplicate keys.

- `--recover`, `-r`

  Do a repair that can fix almost any problem except unique keys that are not unique (which is an extremely unlikely error with `MyISAM` tables). If you want to recover a table, this is the option to try first. You should try `--safe-recover` only if `myisamchk` reports that the table cannot be recovered using `--recover`. (In the unlikely case that `--recover` fails, the data file remains intact.)

  If you have lots of memory, you should increase the value of `myisam_sort_buffer_size`.

- `--safe-recover`, `-o`

  Do a repair using an old recovery method that reads through all rows in order and updates all index trees based on the rows found. This is an order of magnitude slower than `--recover`, but can handle a couple of very unlikely cases that `--recover` cannot. This recovery method also uses much less disk space than `--recover`. Normally, you should repair first using `--recover`, and then with `--safe-recover` only if `--recover` fails.

  If you have lots of memory, you should increase the value of `key_buffer_size`.

- `--set-collation=name`

  Specify the collation to use for sorting table indexes. The character set name is implied by the first part of the collation name.

- `--sort-recover`, `-n`

  Force `myisamchk` to use sorting to resolve the keys even if the temporary files would be very large.

- `--tmpdir=dir_name`, `-t dir_name`

  The path of the directory to be used for storing temporary files. If this is not set, `myisamchk` uses the value of the `TMPDIR` environment variable. `--tmpdir` can be set to a list of directory paths that are used successively in round-robin fashion for creating temporary files. The separator character between directory names is the colon (`:`) on Unix and the semicolon (`;`) on Windows.

- `--unpack`, `-u`

  Unpack a table that was packed with `myisampack`.

4.6.4.4 Other myisamchk Options

`myisamchk` supports the following options for actions other than table checks and repairs:

- `--analyze`, `-a`

  Analyze the distribution of key values. This improves join performance by enabling the join optimizer to better choose the order in which to join the tables and which indexes it should use. To obtain information about the key distribution, use a `myisamchk --description --verbose tbl_name` command or the `SHOW INDEX FROM tbl_name` statement.

- `--block-search=offset`, `-b offset`

  Find the record that a block at the given offset belongs to.

- `--description`, `-d`

  Print some descriptive information about the table. Specifying the `--verbose` option once or twice produces additional information. See Section 4.6.4.5, “Obtaining Table Information with myisamchk”.

- `--set-auto-increment[=value]`, `-A[value]`

  Force `AUTO_INCREMENT` numbering for new records to start at the given value (or higher, if there are existing records with `AUTO_INCREMENT` values this large). If *`value`* is not specified, `AUTO_INCREMENT` numbers for new records begin with the largest value currently in the table, plus one.

- `--sort-index`, `-S`

  Sort the index tree blocks in high-low order. This optimizes seeks and makes table scans that use indexes faster.

- `--sort-records=N`, `-R N`

  Sort records according to a particular index. This makes your data much more localized and may speed up range-based `SELECT` and `ORDER BY` operations that use this index. (The first time you use this option to sort a table, it may be very slow.) To determine a table's index numbers, use `SHOW INDEX`, which displays a table's indexes in the same order that `myisamchk` sees them. Indexes are numbered beginning with 1.

  If keys are not packed (`PACK_KEYS=0`), they have the same length, so when `myisamchk` sorts and moves records, it just overwrites record offsets in the index. If keys are packed (`PACK_KEYS=1`), `myisamchk` must unpack key blocks first, then re-create indexes and pack the key blocks again. (In this case, re-creating indexes is faster than updating offsets for each index.)

4.6.4.5 Obtaining Table Information with myisamchk

To obtain a description of a `MyISAM` table or statistics about it, use the commands shown here. The output from these commands is explained later in this section.

- `myisamchk -d tbl_name`

  Runs `myisamchk` in “describe mode” to produce a description of your table. If you start the MySQL server with external locking disabled, `myisamchk` may report an error for a table that is updated while it runs. However, because `myisamchk` does not change the table in describe mode, there is no risk of destroying data.

- `myisamchk -dv tbl_name`

  Adding `-v` runs `myisamchk` in verbose mode so that it produces more information about the table. Adding `-v` a second time produces even more information.

- `myisamchk -eis tbl_name`

  Shows only the most important information from a table. This operation is slow because it must read the entire table.

- `myisamchk -eiv tbl_name`

  This is like `-eis`, but tells you what is being done.

The *`tbl_name`* argument can be either the name of a `MyISAM` table or the name of its index file, as described in Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”. Multiple *`tbl_name`* arguments can be given.

Suppose that a table named `person` has the following structure. (The `MAX_ROWS` table option is included so that in the example output from `myisamchk` shown later, some values are smaller and fit the output format more easily.)

```
CREATE TABLE person
(
  id         INT NOT NULL AUTO_INCREMENT,
  last_name  VARCHAR(20) NOT NULL,
  first_name VARCHAR(20) NOT NULL,
  birth      DATE,
  death      DATE,
  PRIMARY KEY (id),
  INDEX (last_name, first_name),
  INDEX (birth)
) MAX_ROWS = 1000000 ENGINE=MYISAM;
```

Suppose also that the table has these data and index file sizes:

```
-rw-rw----  1 mysql  mysql  9347072 Aug 19 11:47 person.MYD
-rw-rw----  1 mysql  mysql  6066176 Aug 19 11:47 person.MYI
```

Example of `myisamchk -dvv` output:

```
MyISAM file:         person
Record format:       Packed
Character set:       utf8mb4_0900_ai_ci (255)
File-version:        1
Creation time:       2017-03-30 21:21:30
Status:              checked,analyzed,optimized keys,sorted index pages
Auto increment key:              1  Last value:                306688
Data records:               306688  Deleted blocks:                 0
Datafile parts:             306688  Deleted data:                   0
Datafile pointer (bytes):        4  Keyfile pointer (bytes):        3
Datafile length:           9347072  Keyfile length:           6066176
Max datafile length:    4294967294  Max keyfile length:   17179868159
Recordlength:                   54

table description:
Key Start Len Index   Type                     Rec/key         Root  Blocksize
1   2     4   unique  long                           1                    1024
2   6     80  multip. varchar prefix                 0                    1024
    87    80          varchar                        0
3   168   3   multip. uint24 NULL                    0                    1024

Field Start Length Nullpos Nullbit Type
1     1     1
2     2     4                    no zeros
3     6     81                   varchar
4     87    81                   varchar
5     168   3      1       1     no zeros
6     171   3      1       2     no zeros
```

Explanations for the types of information `myisamchk` produces are given here. “Keyfile” refers to the index file. “Record” and “row” are synonymous, as are “field” and “column.”

The initial part of the table description contains these values:

- `MyISAM file`

  Name of the `MyISAM` (index) file.

- `Record format`

  The format used to store table rows. The preceding examples use `Fixed length`. Other possible values are `Compressed` and `Packed`. (`Packed` corresponds to what `SHOW TABLE STATUS` reports as `Dynamic`.)

- `Chararacter set`

  The table default character set.

- `File-version`

  Version of `MyISAM` format. Always 1.

- `Creation time`

  When the data file was created.

- `Recover time`

  When the index/data file was last reconstructed.

- `Status`

  Table status flags. Possible values are `crashed`, `open`, `changed`, `analyzed`, `optimized keys`, and `sorted index pages`.

- `Auto increment key`, `Last value`

  The key number associated the table's `AUTO_INCREMENT` column, and the most recently generated value for this column. These fields do not appear if there is no such column.

- `Data records`

  The number of rows in the table.

- `Deleted blocks`

  How many deleted blocks still have reserved space. You can optimize your table to minimize this space. See Section 7.6.4, “MyISAM Table Optimization”.

- `Datafile parts`

  For dynamic-row format, this indicates how many data blocks there are. For an optimized table without fragmented rows, this is the same as `Data records`.

- `Deleted data`

  How many bytes of unreclaimed deleted data there are. You can optimize your table to minimize this space. See Section 7.6.4, “MyISAM Table Optimization”.

- `Datafile pointer`

  The size of the data file pointer, in bytes. It is usually 2, 3, 4, or 5 bytes. Most tables manage with 2 bytes, but this cannot be controlled from MySQL yet. For fixed tables, this is a row address. For dynamic tables, this is a byte address.

- `Keyfile pointer`

  The size of the index file pointer, in bytes. It is usually 1, 2, or 3 bytes. Most tables manage with 2 bytes, but this is calculated automatically by MySQL. It is always a block address.

- `Max datafile length`

  How long the table data file can become, in bytes.

- `Max keyfile length`

  How long the table index file can become, in bytes.

- `Recordlength`

How much space each row takes, in bytes.

The `table description` part of the output includes a list of all keys in the table. For each key, `myisamchk` displays some low-level information:

- `Key`

  This key's number. This value is shown only for the first column of the key. If this value is missing, the line corresponds to the second or later column of a multiple-column key. For the table shown in the example, there are two `table description` lines for the second index. This indicates that it is a multiple-part index with two parts.

- `Start`

  Where in the row this portion of the index starts.

- `Len`

  How long this portion of the index is. For packed numbers, this should always be the full length of the column. For strings, it may be shorter than the full length of the indexed column, because you can index a prefix of a string column. The total length of a multiple-part key is the sum of the `Len` values for all key parts.

- `Index`

  Whether a key value can exist multiple times in the index. Possible values are `unique` or `multip.` (multiple).

- `Type`

  What data type this portion of the index has. This is a `MyISAM` data type with the possible values `packed`, `stripped`, or `empty`.

- `Root`

  Address of the root index block.

- `Blocksize`

  The size of each index block. By default this is 1024, but the value may be changed at compile time when MySQL is built from source.

- `Rec/key`

  This is a statistical value used by the optimizer. It tells how many rows there are per value for this index. A unique index always has a value of 1. This may be updated after a table is loaded (or greatly changed) with `myisamchk -a`. If this is not updated at all, a default value of 30 is given.

The last part of the output provides information about each column:

- `Field`

  The column number.

- `Start`

  The byte position of the column within table rows.

- `Length`

  The length of the column in bytes.

- `Nullpos`, `Nullbit`

For columns that can be NULL, MyISAM stores NULL values as a flag in a byte. Depending on how many nullable columns there are, there can be one or more bytes used for this purpose. The Nullpos and Nullbit values, if nonempty, indicate which byte and bit contains that flag indicating whether the column is NULL.

The position and number of bytes used to store NULL flags is shown in the line for field 1. This is why there are six Field lines for the person table even though it has only five columns.

- Type

  The data type. The value may contain any of the following descriptors:

  - constant

    All rows have the same value.

  - no endspace

    Do not store endspace.

  - no endspace, not_always

    Do not store endspace and do not do endspace compression for all values.

  - no endspace, no empty

    Do not store endspace. Do not store empty values.

  - table-lookup

    The column was converted to an ENUM.

  - zerofill(*N*)

    The most significant *N* bytes in the value are always 0 and are not stored.

  - no zeros

    Do not store zeros.

  - always zero

    Zero values are stored using one bit.

- Huff tree

  The number of the Huffman tree associated with the column.

- Bits

  The number of bits used in the Huffman tree.

The Huff tree and Bits fields are displayed if the table has been compressed with myisampack. See Section 4.6.6, “myisampack — Generate Compressed, Read-Only MyISAM Tables”, for an example of this information.

Example of myisamchk -eiv output:

```
Checking MyISAM file: person
Data records:  306688   Deleted blocks:       0
- check file-size
- check record delete-chain
```

```
No recordlinks
- check key delete-chain
block_size 1024:
- check index reference
- check data record references index: 1
Key:  1:  Keyblocks used:  98%  Packed:    0%  Max levels:  3
- check data record references index: 2
Key:  2:  Keyblocks used:  99%  Packed:   97%  Max levels:  3
- check data record references index: 3
Key:  3:  Keyblocks used:  98%  Packed:  -14%  Max levels:  3
Total:    Keyblocks used:  98%  Packed:   89%

- check records and index references
*** LOTS OF ROW NUMBERS DELETED ***

Records:            306688  M.recordlength:       25  Packed:            83%
Recordspace used:       97% Empty space:           2% Blocks/Record:   1.00
Record blocks:      306688  Delete blocks:         0
Record data:       7934464  Deleted data:          0
Lost space:         256512  Linkdata:        1156096

User time 43.08, System time 1.68
Maximum resident set size 0, Integral resident set size 0
Non-physical pagefaults 0, Physical pagefaults 0, Swaps 0
Blocks in 0 out 7, Messages in 0 out 0, Signals 0
Voluntary context switches 0, Involuntary context switches 0
Maximum memory usage: 1046926 bytes (1023k)
```

`myisamchk -eiv` output includes the following information:

- `Data records`

  The number of rows in the table.

- `Deleted blocks`

  How many deleted blocks still have reserved space. You can optimize your table to minimize this space. See Section 7.6.4, “MyISAM Table Optimization”.

- `Key`

  The key number.

- `Keyblocks used`

  What percentage of the keyblocks are used. When a table has just been reorganized with `myisamchk`, the values are very high (very near theoretical maximum).

- `Packed`

  MySQL tries to pack key values that have a common suffix. This can only be used for indexes on `CHAR` and `VARCHAR` columns. For long indexed strings that have similar leftmost parts, this can significantly reduce the space used. In the preceding example, the second key is 40 bytes long and a 97% reduction in space is achieved.

- `Max levels`

  How deep the B-tree for this key is. Large tables with long key values get high values.

- `Records`

  How many rows are in the table.

- `M.recordlength`

  The average row length. This is the exact row length for tables with fixed-length rows, because all rows have the same length.

- `Packed`

  MySQL strips spaces from the end of strings. The `Packed` value indicates the percentage of savings achieved by doing this.

- `Recordspace used`

  What percentage of the data file is used.

- `Empty space`

  What percentage of the data file is unused.

- `Blocks/Record`

  Average number of blocks per row (that is, how many links a fragmented row is composed of). This is always 1.0 for fixed-format tables. This value should stay as close to 1.0 as possible. If it gets too large, you can reorganize the table. See Section 7.6.4, “MyISAM Table Optimization”.

- `Recordblocks`

  How many blocks (links) are used. For fixed-format tables, this is the same as the number of rows.

- `Deleteblocks`

  How many blocks (links) are deleted.

- `Recorddata`

  How many bytes in the data file are used.

- `Deleted data`

  How many bytes in the data file are deleted (unused).

- `Lost space`

  If a row is updated to a shorter length, some space is lost. This is the sum of all such losses, in bytes.

- `Linkdata`

  When the dynamic table format is used, row fragments are linked with pointers (4 to 7 bytes each). `Linkdata` is the sum of the amount of storage used by all such pointers.

4.6.4.6 myisamchk Memory Usage

Memory allocation is important when you run `myisamchk`. `myisamchk` uses no more memory than its memory-related variables are set to. If you are going to use `myisamchk` on very large tables, you should first decide how much memory you want it to use. The default is to use only about 3MB to perform repairs. By using larger values, you can get `myisamchk` to operate faster. For example, if you have more than 512MB RAM available, you could use options such as these (in addition to any other options you might specify):

```
shell> myisamchk --myisam_sort_buffer_size=256M \
           --key_buffer_size=512M \
           --read_buffer_size=64M \
           --write_buffer_size=64M ...
```

Using `--myisam_sort_buffer_size=16M` is probably enough for most cases.

Be aware that `myisamchk` uses temporary files in `TMPDIR`. If `TMPDIR` points to a memory file system, out of memory errors can easily occur. If this happens, run `myisamchk` with the `--tmpdir=dir_name` option to specify a directory located on a file system that has more space.

When performing repair operations, myisamchk also needs a lot of disk space:

- Twice the size of the data file (the original file and a copy). This space is not needed if you do a repair with --quick; in this case, only the index file is re-created. *This space must be available on the same file system as the original data file*, as the copy is created in the same directory as the original.
- Space for the new index file that replaces the old one. The old index file is truncated at the start of the repair operation, so you usually ignore this space. This space must be available on the same file system as the original data file.
- When using --recover or --sort-recover (but not when using --safe-recover), you need space on disk for sorting. This space is allocated in the temporary directory (specified by TMPDIR or --tmpdir=*dir_name*). The following formula yields the amount of space required:

```
(largest_key + row_pointer_length) * number_of_rows * 2
```

  You can check the length of the keys and the *row_pointer_length* with myisamchk -dv *tbl_name* (see Section 4.6.4.5, "Obtaining Table Information with myisamchk"). The *row_pointer_length* and *number_of_rows* values are the Datafile pointer and Data records values in the table description. To determine the *largest_key* value, check the Key lines in the table description. The Len column indicates the number of bytes for each key part. For a multiple-column index, the key size is the sum of the Len values for all key parts.

If you have a problem with disk space during repair, you can try --safe-recover instead of --recover.

4.6.5 myisamlog — Display MyISAM Log File Contents

myisamlog processes the contents of a MyISAM log file. To create such a file, start the server with a --log-isam=log_file option.

Invoke myisamlog like this:

```
shell> myisamlog [options] [file_name [tbl_name] ...]
```

The default operation is update (-u). If a recovery is done (-r), all writes and possibly updates and deletes are done and errors are only counted. The default log file name is myisam.log if no *log_file* argument is given. If tables are named on the command line, only those tables are updated.

myisamlog supports the following options:

- -?, -I

  Display a help message and exit.

- -c *N*

  Execute only *N* commands.

- -f *N*

  Specify the maximum number of open files.

- -i

  Display extra information before exiting.

- -o *offset*

  Specify the starting offset.

- `-p N`

  Remove `N` components from path.

- `-r`

  Perform a recovery operation.

- `-R record_pos_file record_pos`

  Specify record position file and record position.

- `-u`

  Perform an update operation.

- `-v`

  Verbose mode. Print more output about what the program does. This option can be given multiple times to produce more and more output.

- `-w write_file`

  Specify the write file.

- `-V`

  Display version information.

4.6.6 myisampack — Generate Compressed, Read-Only MyISAM Tables

The `myisampack` utility compresses `MyISAM` tables. `myisampack` works by compressing each column in the table separately. Usually, `myisampack` packs the data file 40% to 70%.

When the table is used later, the server reads into memory the information needed to decompress columns. This results in much better performance when accessing individual rows, because you only have to uncompress exactly one row.

MySQL uses `mmap()` when possible to perform memory mapping on compressed tables. If `mmap()` does not work, MySQL falls back to normal read/write file operations.

Please note the following:

- If the `mysqld` server was invoked with external locking disabled, it is not a good idea to invoke `myisampack` if the table might be updated by the server during the packing process. It is safest to compress tables with the server stopped.
- After packing a table, it becomes read only. This is generally intended (such as when accessing packed tables on a CD).
- `myisampack` does not support partitioned tables.

Invoke `myisampack` like this:

```
shell> myisampack [options] file_name ...
```

Each file name argument should be the name of an index (`.MYI`) file. If you are not in the database directory, you should specify the path name to the file. It is permissible to omit the `.MYI` extension.

After you compress a table with `myisampack`, use `myisamchk -rq` to rebuild its indexes. Section 4.6.4, “myisamchk — MyISAM Table-Maintenance Utility”.

myisampack supports the following options. It also reads option files and supports the options for processing them described at Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- --help, -?

  Display a help message and exit.

- --backup, -b

  Make a backup of each table's data file using the name *tbl_name*.OLD.

- --character-sets-dir=*dir_name*

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- --debug[=*debug_options*], -# [*debug_options*]

  Write a debugging log. A typical *debug_options* string is d:t:o,*file_name*. The default is d:t:o.

- --force, -f

  Produce a packed table even if it becomes larger than the original or if the intermediate file from an earlier invocation of myisampack exists. (myisampack creates an intermediate file named *tbl_name*.TMD in the database directory while it compresses the table. If you kill myisampack, the .TMD file might not be deleted.) Normally, myisampack exits with an error if it finds that *tbl_name*.TMD exists. With --force, myisampack packs the table anyway.

- --join=*big_tbl_name*, -j *big_tbl_name*

  Join all tables named on the command line into a single packed table *big_tbl_name*. All tables that are to be combined *must* have identical structure (same column names and types, same indexes, and so forth).

  big_tbl_name must not exist prior to the join operation. All source tables named on the command line to be merged into *big_tbl_name* must exist. The source tables are read for the join operation but not modified.

- --silent, -s

  Silent mode. Write output only when errors occur.

- --test, -t

  Do not actually pack the table, just test packing it.

- --tmpdir=*dir_name*, -T *dir_name*

  Use the named directory as the location where myisampack creates temporary files.

- --verbose, -v

  Verbose mode. Write information about the progress of the packing operation and its result.

- --version, -V

  Display version information and exit.

- --wait, -w

  Wait and retry if the table is in use. If the mysqld server was invoked with external locking disabled, it is not a good idea to invoke myisampack if the table might be updated by the server during the packing process.

The following sequence of commands illustrates a typical table compression session:

```
shell> ls -l station.*
-rw-rw-r--   1 monty    my         994128 Apr 17 19:00 station.MYD
-rw-rw-r--   1 monty    my          53248 Apr 17 19:00 station.MYI

shell> myisamchk -dvv station

MyISAM file:     station
Isam-version:  2
Creation time: 1996-03-13 10:08:58
Recover time:  1997-02-02  3:06:43
Data records:              1192  Deleted blocks:              0
Datafile parts:            1192  Deleted data:                0
Datafile pointer (bytes):     2  Keyfile pointer (bytes):     2
Max datafile length:   54657023  Max keyfile length:   33554431
Recordlength:               834
Record format: Fixed length

table description:
Key Start Len Index   Type                 Root  Blocksize    Rec/key
1   2     4   unique  unsigned long        1024       1024          1
2   32    30  multip. text                10240       1024          1

Field Start Length Type
1     1     1
2     2     4
3     6     4
4     10    1
5     11    20
6     31    1
7     32    30
8     62    35
9     97    35
10    132   35
11    167   4
12    171   16
13    187   35
14    222   4
15    226   16
16    242   20
17    262   20
18    282   20
19    302   30
20    332   4
21    336   4
22    340   1
23    341   8
24    349   8
25    357   8
26    365   2
27    367   2
28    369   4
29    373   4
30    377   1
31    378   2
32    380   8
33    388   4
34    392   4
35    396   4
36    400   4
37    404   1
38    405   4
39    409   4
40    413   4
41    417   4
42    421   4
43    425   4
44    429   20
45    449   30
```

```
46    479   1
47    480   1
48    481   79
49    560   79
50    639   79
51    718   79
52    797   8
53    805   1
54    806   1
55    807   20
56    827   4
57    831   4

shell> myisampack station.MYI
Compressing station.MYI: (1192 records)
- Calculating statistics

normal:     20  empty-space:   16  empty-zero:     12  empty-fill:  11
pre-space:   0  end-space:     12  table-lookups:   5  zero:         7
Original trees:  57  After join: 17
- Compressing file
87.14%
Remember to run myisamchk -rq on compressed tables

shell> myisamchk -rq station
- check record delete-chain
- recovering (with sort) MyISAM-table 'station'
Data records: 1192
- Fixing index 1
- Fixing index 2

shell> mysqladmin -uroot flush-tables

shell> ls -l station.*
-rw-rw-r--   1 monty    my         127874 Apr 17 19:00 station.MYD
-rw-rw-r--   1 monty    my          55296 Apr 17 19:04 station.MYI

shell> myisamchk -dvv station

MyISAM file:     station
Isam-version:  2
Creation time: 1996-03-13 10:08:58
Recover time:  1997-04-17 19:04:26
Data records:               1192  Deleted blocks:              0
Datafile parts:             1192  Deleted data:                0
Datafile pointer (bytes):      3  Keyfile pointer (bytes):     1
Max datafile length:    16777215  Max keyfile length:     131071
Recordlength:               834
Record format: Compressed

table description:
Key Start Len Index   Type                 Root  Blocksize    Rec/key
1   2     4   unique  unsigned long       10240       1024          1
2   32    30  multip. text                54272       1024          1

Field Start Length Type                          Huff tree  Bits
1     1     1      constant                            1     0
2     2     4      zerofill(1)                         2     9
3     6     4      no zeros, zerofill(1)               2     9
4     10    1                                          3     9
5     11    20     table-lookup                        4     0
6     31    1                                          3     9
7     32    30     no endspace, not_always             5     9
8     62    35     no endspace, not_always, no empty   6     9
9     97    35     no empty                            7     9
10    132   35     no endspace, not_always, no empty   6     9
11    167   4      zerofill(1)                         2     9
12    171   16     no endspace, not_always, no empty   5     9
13    187   35     no endspace, not_always, no empty   6     9
14    222   4      zerofill(1)                         2     9
15    226   16     no endspace, not_always, no empty   5     9
16    242   20     no endspace, not_always             8     9
```

```
17    262   20     no endspace, no empty               8     9
18    282   20     no endspace, no empty               5     9
19    302   30     no endspace, no empty               6     9
20    332   4      always zero                         2     9
21    336   4      always zero                         2     9
22    340   1                                          3     9
23    341   8      table-lookup                        9     0
24    349   8      table-lookup                       10     0
25    357   8      always zero                         2     9
26    365   2                                          2     9
27    367   2      no zeros, zerofill(1)               2     9
28    369   4      no zeros, zerofill(1)               2     9
29    373   4      table-lookup                       11     0
30    377   1                                          3     9
31    378   2      no zeros, zerofill(1)               2     9
32    380   8      no zeros                            2     9
33    388   4      always zero                         2     9
34    392   4      table-lookup                       12     0
35    396   4      no zeros, zerofill(1)              13     9
36    400   4      no zeros, zerofill(1)               2     9
37    404   1                                          2     9
38    405   4      no zeros                            2     9
39    409   4      always zero                         2     9
40    413   4      no zeros                            2     9
41    417   4      always zero                         2     9
42    421   4      no zeros                            2     9
43    425   4      always zero                         2     9
44    429   20     no empty                            3     9
45    449   30     no empty                            3     9
46    479   1                                         14     4
47    480   1                                         14     4
48    481   79     no endspace, no empty              15     9
49    560   79     no empty                            2     9
50    639   79     no empty                            2     9
51    718   79     no endspace                        16     9
52    797   8      no empty                            2     9
53    805   1                                         17     1
54    806   1                                          3     9
55    807   20     no empty                            3     9
56    827   4      no zeros, zerofill(2)               2     9
57    831   4      no zeros, zerofill(1)               2     9
```

`myisampack` displays the following kinds of information:

- `normal`

  The number of columns for which no extra packing is used.

- `empty-space`

  The number of columns containing values that are only spaces. These occupy one bit.

- `empty-zero`

  The number of columns containing values that are only binary zeros. These occupy one bit.

- `empty-fill`

  The number of integer columns that do not occupy the full byte range of their type. These are changed to a smaller type. For example, a `BIGINT` column (eight bytes) can be stored as a `TINYINT` column (one byte) if all its values are in the range from `-128` to `127`.

- `pre-space`

  The number of decimal columns that are stored with leading spaces. In this case, each value contains a count for the number of leading spaces.

- `end-space`

The number of columns that have a lot of trailing spaces. In this case, each value contains a count for the number of trailing spaces.

- `table-lookup`

  The column had only a small number of different values, which were converted to an `ENUM` before Huffman compression.

- `zero`

  The number of columns for which all values are zero.

- `Original trees`

  The initial number of Huffman trees.

- `After join`

  The number of distinct Huffman trees left after joining trees to save some header space.

After a table has been compressed, the `Field` lines displayed by `myisamchk -dvv` include additional information about each column:

- `Type`

  The data type. The value may contain any of the following descriptors:

  - `constant`

    All rows have the same value.

  - `no endspace`

    Do not store endspace.

  - `no endspace, not_always`

    Do not store endspace and do not do endspace compression for all values.

  - `no endspace, no empty`

    Do not store endspace. Do not store empty values.

  - `table-lookup`

    The column was converted to an `ENUM`.

  - `zerofill(N)`

    The most significant *`N`* bytes in the value are always 0 and are not stored.

  - `no zeros`

    Do not store zeros.

  - `always zero`

    Zero values are stored using one bit.

- `Huff tree`

  The number of the Huffman tree associated with the column.

- `Bits`

  The number of bits used in the Huffman tree.

After you run `myisampack`, use `myisamchk` to re-create any indexes. At this time, you can also sort the index blocks and create statistics needed for the MySQL optimizer to work more efficiently:

```
shell> myisamchk -rq --sort-index --analyze tbl_name.MYI
```

After you have installed the packed table into the MySQL database directory, you should execute `mysqladmin flush-tables` to force `mysqld` to start using the new table.

To unpack a packed table, use the `--unpack` option to `myisamchk`.

4.6.7 `mysql_config_editor` — MySQL Configuration Utility

The `mysql_config_editor` utility enables you to store authentication credentials in an encrypted login path file named `.mylogin.cnf`. The file location is the `%APPDATA%\MySQL` directory on Windows and the current user's home directory on non-Windows systems. The file can be read later by MySQL client programs to obtain authentication credentials for connecting to MySQL Server.

The unencrypted format of the `.mylogin.cnf` login path file consists of option groups, similar to other option files. Each option group in `.mylogin.cnf` is called a “login path,” which is a group that permits only certain options: `host`, `user`, `password`, `port` and `socket`. Think of a login path option group as a set of options that specify which MySQL server to connect to and which account to authenticate as. Here is an unencrypted example:

```
[client]
user = mydefaultname
password = mydefaultpass
host = 127.0.0.1
[mypath]
user = myothername
password = myotherpass
host = localhost
```

When you invoke a client program to connect to the server, the client uses `.mylogin.cnf` in conjunction with other option files. Its precedence is higher than other option files, but less than options specified explicitly on the client command line. For information about the order in which option files are used, see Section 4.2.6, “Using Option Files”.

To specify an alternate login path file name, set the `MYSQL_TEST_LOGIN_FILE` environment variable. This variable is recognized by `mysql_config_editor`, by standard MySQL clients (`mysql`, `mysqladmin`, and so forth), and by the `mysql-test-run.pl` testing utility.

Programs use groups in the login path file as follows:

- `mysql_config_editor` operates on the `client` login path by default if you specify no `--login-path=`*name* option to indicate explicitly which login path to use.
- Without a `--login-path` option, client programs read the same option groups from the login path file that they read from other option files. Consider this command:

  ```
  shell> mysql
  ```

  By default, the `mysql` client reads the `[client]` and `[mysql]` groups from other option files, so it reads them from the login path file as well.
- With a `--login-path` option, client programs additionally read the named login path from the login path file. The option groups read from other option files remain the same. Consider this command:

```
shell> mysql --login-path=mypath
```

The mysql client reads [client] and [mysql] from other option files, and [client], [mysql], and [mypath] from the login path file.

- Client programs read the login path file even when the --no-defaults option is used. This permits passwords to be specified in a safer way than on the command line even if --no-defaults is present.

mysql_config_editor encrypts the .mylogin.cnf file so it cannot be read as cleartext, and its contents when decrypted by client programs are used only in memory. In this way, passwords can be stored in a file in non-cleartext format and used later without ever needing to be exposed on the command line or in an environment variable. mysql_config_editor provides a print command for displaying the login path file contents, but even in this case, password values are masked so as never to appear in a way that other users can see them.

The encryption used by mysql_config_editor prevents passwords from appearing in .mylogin.cnf as cleartext and provides a measure of security by preventing inadvertent password exposure. For example, if you display a regular unencrypted my.cnf option file on the screen, any passwords it contains are visible for anyone to see. With .mylogin.cnf, that is not true. But the encryption used will not deter a determined attacker and you should not consider it unbreakable. A user who can gain system administration privileges on your machine to access your files could decrypt the .mylogin.cnf file with some effort.

The login path file must be readable and writable to the current user, and inaccessible to other users. Otherwise, mysql_config_editor ignores it, and client programs do not use it, either.

Invoke mysql_config_editor like this:

```
shell> mysql_config_editor [program_options] command [command_options]
```

If the login path file does not exist, mysql_config_editor creates it.

Command arguments are given as follows:

- *program_options* consists of general mysql_config_editor options.
- command indicates what action to perform on the .mylogin.cnf login path file. For example, set writes a login path to the file, remove removes a login path, and print displays login path contents.
- *command_options* indicates any additional options specific to the command, such as the login path name and the values to use in the login path.

The position of the command name within the set of program arguments is significant. For example, these command lines have the same arguments, but produce different results:

```
shell> mysql_config_editor --help set
shell> mysql_config_editor set --help
```

The first command line displays a general mysql_config_editor help message, and ignores the set command. The second command line displays a help message specific to the set command.

Suppose that you want to establish a client login path that defines your default connection parameters, and an additional login path named remote for connecting to the MySQL server the host remote.example.com. You want to log in as follows:

- By default, to the local server with a user name and password of localuser and localpass
- To the remote server with a user name and password of remoteuser and remotepass

To set up the login paths in the .mylogin.cnf file, use the following set commands. Enter each command on a single line, and enter the appropriate passwords when prompted:

```
shell> mysql_config_editor set --login-path=client
         --host=localhost --user=localuser --password
Enter password: enter password "localpass" here
shell> mysql_config_editor set --login-path=remote
         --host=remote.example.com --user=remoteuser --password
Enter password: enter password "remotepass" here
```

mysql_config_editor uses the client login path by default, so the --login-path=client option can be omitted from the first command without changing its effect.

To see what mysql_config_editor writes to the .mylogin.cnf file, use the print command:

```
shell> mysql_config_editor print --all
[client]
user = localuser
password = *****
host = localhost
[remote]
user = remoteuser
password = *****
host = remote.example.com
```

The print command displays each login path as a set of lines beginning with a group header indicating the login path name in square brackets, followed by the option values for the login path. Password values are masked and do not appear as cleartext.

If you do not specify --all to display all login paths or --login-path=*name* to display a named login path, the print command displays the client login path by default, if there is one.

As shown by the preceding example, the login path file can contain multiple login paths. In this way, mysql_config_editor makes it easy to set up multiple “personalities” for connecting to different MySQL servers, or for connecting to a given server using different accounts. Any of these can be selected by name later using the --login-path option when you invoke a client program. For example, to connect to the remote server, use this command:

```
shell> mysql --login-path=remote
```

Here, mysql reads the [client] and [mysql] option groups from other option files, and the [client], [mysql], and [remote] groups from the login path file.

To connect to the local server, use this command:

```
shell> mysql --login-path=client
```

Because mysql reads the client and mysql login paths by default, the --login-path option does not add anything in this case. That command is equivalent to this one:

```
shell> mysql
```

Options read from the login path file take precedence over options read from other option files. Options read from login path groups appearing later in the login path file take precedence over options read from groups appearing earlier in the file.

mysql_config_editor adds login paths to the login path file in the order you create them, so you should create more general login paths first and more specific paths later. If you need to move a login path within the file, you can remove it, then recreate it to add it to the end. For example, a client login path is more general because it is read by all client programs, whereas a mysqldump login path is read only by mysqldump. Options specified later override options specified earlier, so putting the

login paths in the order client, mysqldump enables mysqldump-specific options to override client options.

When you use the set command with mysql_config_editor to create a login path, you need not specify all possible option values (host name, user name, password, port, socket). Only those values given are written to the path. Any missing values required later can be specified when you invoke a client path to connect to the MySQL server, either in other option files or on the command line. Any options specified on the command line override those specified in the login path file or other option files. For example, if the credentials in the remote login path also apply for the host remote2.example.com, connect to the server on that host like this:

```
shell> mysql --login-path=remote --host=remote2.example.com
```

mysql_config_editor General Options

mysql_config_editor supports the following general options, which may be used preceding any command named on the command line. For descriptions of command-specific options, see mysql_config_editor Commands and Command-Specific Options.

Table 4.17 mysql_config_editor General Options

Format	Description
--debug	Write debugging log
--help	Display help message and exit
--verbose	Verbose mode
--version	Display version information and exit

- --help, -?

  Display a general help message and exit.

  To see a command-specific help message, invoke mysql_config_editor as follows, where *command* is a command other than help:

  ```
  shell> mysql_config_editor command --help
  ```

- --debug[=*debug_options*], -# *debug_options*

  Write a debugging log. A typical *debug_options* string is d:t:o,*file_name*. The default is d:t:o,/tmp/mysql_config_editor.trace.

- --verbose, -v

  Verbose mode. Print more information about what the program does. This option may be helpful in diagnosing problems if an operation does not have the effect you expect.

- --version, -V

  Display version information and exit.

mysql_config_editor Commands and Command-Specific Options

This section describes the permitted mysql_config_editor commands, and, for each one, the command-specific options permitted following the command name on the command line.

In addition, mysql_config_editor supports general options that can be used preceding any command. For descriptions of these options, see mysql_config_editor General Options.

mysql_config_editor supports these commands:

- `help`

  Display a general help message and exit. This command takes no following options.

  To see a command-specific help message, invoke `mysql_config_editor` as follows, where *`command`* is a command other than `help`:

```
shell> mysql_config_editor command --help
```

- `print [options]`

  Print the contents of the login path file in unencrypted form, with the exception that passwords are displayed as `*****`.

  The default login path name is `client` if no login path is named. If both `--all` and `--login-path` are given, `--all` takes precedence.

  The `print` command permits these options following the command name:

  - `--help`, `-?`

    Display a help message for the `print` command and exit.

    To see a general help message, use `mysql_config_editor --help`.

  - `--all`

    Print the contents of all login paths in the login path file.

  - `--login-path=name`, `-G name`

    Print the contents of the named login path.

- `remove [options]`

  Remove a login path from the login path file, or modify a login path by removing options from it.

  This command removes from the login path only such options as are specified with the `--host`, `--password`, `--port`, `--socket`, and `--user` options. If none of those options are given, `remove` removes the entire login path. For example, this command removes only the `user` option from the `mypath` login path rather than the entire `mypath` login path:

```
shell> mysql_config_editor remove --login-path=mypath --user
```

  This command removes the entire `mypath` login path:

```
shell> mysql_config_editor remove --login-path=mypath
```

  The `remove` command permits these options following the command name:

  - `--help`, `-?`

    Display a help message for the `remove` command and exit.

    To see a general help message, use `mysql_config_editor --help`.

  - `--host`, `-h`

    Remove the host name from the login path.

  - `--login-path=name`, `-G name`

The login path to remove or modify. The default login path name is `client` if this option is not given.

- `--password`, `-p`

  Remove the password from the login path.

- `--port`, `-P`

  Remove the TCP/IP port number from the login path.

- `--socket`, `-S`

  Remove the Unix socket file name from the login path.

- `--user`, `-u`

  Remove the user name from the login path.

- `--warn`, `-w`

  Warn and prompt the user for confirmation if the command attempts to remove the default login path (`client`) and `--login-path=client` was not specified. This option is enabled by default; use `--skip-warn` to disable it.

- `reset [options]`

  Empty the contents of the login path file.

  The `reset` command permits these options following the command name:

  - `--help`, `-?`

    Display a help message for the `reset` command and exit.

    To see a general help message, use `mysql_config_editor --help`.

- `set [options]`

  Write a login path to the login path file.

  This command writes to the login path only such options as are specified with the `--host`, `--password`, `--port`, `--socket`, and `--user` options. If none of those options are given, `mysql_config_editor` writes the login path as an empty group.

  The `set` command permits these options following the command name:

  - `--help`, `-?`

    Display a help message for the `set` command and exit.

    To see a general help message, use `mysql_config_editor --help`.

  - `--host=host_name`, `-h host_name`

    The host name to write to the login path.

  - `--login-path=name`, `-G name`

    The login path to create. The default login path name is `client` if this option is not given.

  - `--password`, `-p`

Prompt for a password to write to the login path. After mysql_config_editor displays the prompt, type the password and press Enter. To prevent other users from seeing the password, mysql_config_editor does not echo it.

To specify an empty password, press Enter at the password prompt. The resulting login path written to the login path file will include a line like this:

```
password =
```

- --port=*port_num*, -P *port_num*

  The TCP/IP port number to write to the login path.

- --socket=*file_name*, -S *file_name*

  The Unix socket file name to write to the login path.

- --user=*user_name*, -u *user_name*

  The user name to write to the login path.

- --warn, -w

  Warn and prompt the user for confirmation if the command attempts to overwrite an existing login path. This option is enabled by default; use --skip-warn to disable it.

4.6.8 mysqlbinlog — Utility for Processing Binary Log Files

The server's binary log consists of files containing “events” that describe modifications to database contents. The server writes these files in binary format. To display their contents in text format, use the mysqlbinlog utility. You can also use mysqlbinlog to display the contents of relay log files written by a slave server in a replication setup because relay logs have the same format as binary logs. The binary log and relay log are discussed further in Section 5.4.4, “The Binary Log”, and Section 18.2.4, “Replication Relay and Status Logs”.

Invoke mysqlbinlog like this:

```
shell> mysqlbinlog [options] log_file ...
```

For example, to display the contents of the binary log file named binlog.000003, use this command:

```
shell> mysqlbinlog binlog.0000003
```

The output includes events contained in binlog.000003. For statement-based logging, event information includes the SQL statement, the ID of the server on which it was executed, the timestamp when the statement was executed, how much time it took, and so forth. For row-based logging, the event indicates a row change rather than an SQL statement. See Section 18.2.1, “Replication Formats”, for information about logging modes.

Events are preceded by header comments that provide additional information. For example:

```
# at 141
#100309  9:28:36 server id 123  end_log_pos 245
  Query thread_id=3350  exec_time=11  error_code=0
```

In the first line, the number following at indicates the file offset, or starting position, of the event in the binary log file.

The second line starts with a date and time indicating when the statement started on the server where the event originated. For replication, this timestamp is propagated to slave servers. server id is

the server_id value of the server where the event originated. end_log_pos indicates where the next event starts (that is, it is the end position of the current event + 1). thread_id indicates which thread executed the event. exec_time is the time spent executing the event, on a master server. On a slave, it is the difference of the end execution time on the slave minus the beginning execution time on the master. The difference serves as an indicator of how much replication lags behind the master. error_code indicates the result from executing the event. Zero means that no error occurred.

Note

When using event groups, the file offsets of events may be grouped together and the comments of events may be grouped together. Do not mistake these grouped events for blank file offsets.

The output from mysqlbinlog can be re-executed (for example, by using it as input to mysql) to redo the statements in the log. This is useful for recovery operations after a server crash. For other usage examples, see the discussion later in this section and in Section 7.5, "Point-in-Time (Incremental) Recovery Using the Binary Log".

Normally, you use mysqlbinlog to read binary log files directly and apply them to the local MySQL server. It is also possible to read binary logs from a remote server by using the --read-from-remote-server option. To read remote binary logs, the connection parameter options can be given to indicate how to connect to the server. These options are --host, --password, --port, --protocol, --socket, and --user; they are ignored except when you also use the --read-from-remote-server option.

When running mysqlbinlog against a large binary log, be careful that the filesystem has enough space for the resulting files. To configure the directory that mysqlbinlog uses for temporary files, use the TMPDIR environment variable.

mysqlbinlog supports the following options, which can be specified on the command line or in the [mysqlbinlog] and [client] groups of an option file. For information about option files used by MySQL programs, see Section 4.2.6, "Using Option Files".

Table 4.18 mysqlbinlog Options

Format	Description
--base64-output	Print binary log entries using base-64 encoding
--bind-address	Use specified network interface to connect to MySQL Server
--binlog-row-event-max-size	Binary log max event size
--character-sets-dir	Directory where character sets are installed
--connection-server-id	Used for testing and debugging. See text for applicable default values and other particulars.
--database	List entries for just this database
--debug	Write debugging log
--debug-check	Print debugging information when program exits
--debug-info	Print debugging information, memory, and CPU statistics when program exits
--default-auth	Authentication plugin to use
--defaults-extra-file	Read named option file in addition to usual option files
--defaults-file	Read only named option file
--defaults-group-suffix	Option group suffix value
--disable-log-bin	Disable binary logging
--exclude-gtids	Do not show any of the groups in the GTID set provided
--force-if-open	Read binary log files even if open or not closed properly

Format	Description
--force-read	If mysqlbinlog reads a binary log event that it does not recognize, it prints a warning
--help	Display help message and exit
--hexdump	Display a hex dump of the log in comments
--host	Connect to MySQL server on given host
--idempotent	Cause the server to use idempotent mode while processing binary log updates from this session only
--include-gtids	Show only the groups in the GTID set provided
--local-load	Prepare local temporary files for LOAD DATA INFILE in the specified directory
--login-path	Read login path options from .mylogin.cnf
--no-defaults	Read no option files
--offset	Skip the first N entries in the log
--password	Password to use when connecting to server
--plugin-dir	Directory where plugins are installed
--port	TCP/IP port number to use for connection
--print-defaults	Print default options
--print-table-metadata	Print table metadata
--protocol	Connection protocol to use
--raw	Write events in raw (binary) format to output files
--read-from-remote-master	Read the binary log from a MySQL master rather than reading a local log file
--read-from-remote-server	Read binary log from MySQL server rather than local log file
--result-file	Direct output to named file
--rewrite-db	Create rewrite rules for databases when playing back from logs written in row-based format. Can be used multiple times.
--secure-auth	Do not send passwords to server in old (pre-4.1) format
--server-id	Extract only those events created by the server having the given server ID
--server-id-bits	Tell mysqlbinlog how to interpret server IDs in binary log when log was written by a mysqld having its server-id-bits set to less than the maximum; supported only by MySQL Cluster version of mysqlbinlog
--set-charset	Add a SET NAMES charset_name statement to the output
--shared-memory-base-name	The name of shared memory to use for shared-memory connections
--short-form	Display only the statements contained in the log
--skip-gtids	Do not print any GTIDs; use this when writing a dump file from binary logs containing GTIDs.
--socket	For connections to localhost, the Unix socket file to use
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format

Format	Description
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--start-datetime	Read binary log from first event with timestamp equal to or later than datetime argument
--start-position	Read binary log from first event with position equal to or greater than argument
--stop-datetime	Stop reading binary log at first event with timestamp equal to or greater than datetime argument
--stop-never	Stay connected to server after reading last binary log file
--stop-never-slave-server-id	Slave server ID to report when connecting to server
--stop-position	Stop reading binary log at first event with position equal to or greater than argument
--tls-version	Protocols permitted for secure connections
--to-last-log	Do not stop at the end of requested binary log from a MySQL server, but rather continue printing to end of last binary log
--user	MySQL user name to use when connecting to server
--verbose	Reconstruct row events as SQL statements
--verify-binlog-checksum	Verify checksums in binary log
--version	Display version information and exit

- `--help`, `-?`

  Display a help message and exit.

- `--base64-output=value`

  This option determines when events should be displayed encoded as base-64 strings using `BINLOG` statements. The option has these permissible values (not case sensitive):

  - `AUTO` ("automatic") or `UNSPEC` ("unspecified") displays `BINLOG` statements automatically when necessary (that is, for format description events and row events). If no `--base64-output` option is given, the effect is the same as `--base64-output=AUTO`.

    Note

    Automatic `BINLOG` display is the only safe behavior if you intend to use the output of `mysqlbinlog` to re-execute binary log file contents. The other option values are intended only for debugging or testing purposes because they may produce output that does not include all events in executable form.

  - `NEVER` causes `BINLOG` statements not to be displayed. `mysqlbinlog` exits with an error if a row event is found that must be displayed using `BINLOG`.

  - `DECODE-ROWS` specifies to `mysqlbinlog` that you intend for row events to be decoded and displayed as commented SQL statements by also specifying the `--verbose` option. Like `NEVER`, `DECODE-ROWS` suppresses display of `BINLOG` statements, but unlike `NEVER`, it does not exit with an error if a row event is found.

For examples that show the effect of `--base64-output` and `--verbose` on row event output, see Section 4.6.8.2, “mysqlbinlog Row Event Display”.

- `--bind-address=ip_address`

  On a computer having multiple network interfaces, use this option to select which interface to use for connecting to the MySQL server.

- `--binlog-row-event-max-size=N`

Command-Line Format	`--binlog-row-event-max-size=#`	
Permitted Values (64-bit platforms)	**Type**	numeric
	Default	`4294967040`
	Min Value	`256`
	Max Value	`18446744073709547520`

  Specify the maximum size of a row-based binary log event, in bytes. Rows are grouped into events smaller than this size if possible. The value should be a multiple of 256. The default is 4GB.

- `--character-sets-dir=dir_name`

  The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--connection-server-id=server_id`

  This option is used to test a MySQL server for support of the `BINLOG_DUMP_NON_BLOCK` connection flag, which was inadvertently removed in MySQL 5.6.5, and restored in MySQL 5.7.5 (Bug #18000079, Bug #71178). It is not required for normal operations.

  The effective default and minimum values for this option depend on whether `mysqlbinlog` is run in blocking mode or non-blocking mode. When `mysqlbinlog` is run in blocking mode, the default (and minimum) value is 1; when run in non-blocking mode, the default (and minimum) value is 0.

- `--database=db_name`, `-d db_name`

  This option causes `mysqlbinlog` to output entries from the binary log (local log only) that occur while *`db_name`* is been selected as the default database by `USE`.

  The `--database` option for `mysqlbinlog` is similar to the `--binlog-do-db` option for `mysqld`, but can be used to specify only one database. If `--database` is given multiple times, only the last instance is used.

  The effects of this option depend on whether the statement-based or row-based logging format is in use, in the same way that the effects of `--binlog-do-db` depend on whether statement-based or row-based logging is in use.

  Statement-based logging. The `--database` option works as follows:

  - While *`db_name`* is the default database, statements are output whether they modify tables in *`db_name`* or a different database.
  - Unless *`db_name`* is selected as the default database, statements are not output, even if they modify tables in *`db_name`*.
  - There is an exception for `CREATE DATABASE`, `ALTER DATABASE`, and `DROP DATABASE`. The database being *created, altered, or dropped* is considered to be the default database when determining whether to output the statement.

Suppose that the binary log was created by executing these statements using statement-based-logging:

```
INSERT INTO test.t1 (i) VALUES(100);
INSERT INTO db2.t2 (j)  VALUES(200);
USE test;
INSERT INTO test.t1 (i) VALUES(101);
INSERT INTO t1 (i)      VALUES(102);
INSERT INTO db2.t2 (j)  VALUES(201);
USE db2;
INSERT INTO test.t1 (i) VALUES(103);
INSERT INTO db2.t2 (j)  VALUES(202);
INSERT INTO t2 (j)      VALUES(203);
```

`mysqlbinlog --database=test` does not output the first two `INSERT` statements because there is no default database. It outputs the three `INSERT` statements following `USE test`, but not the three `INSERT` statements following `USE db2`.

`mysqlbinlog --database=db2` does not output the first two `INSERT` statements because there is no default database. It does not output the three `INSERT` statements following `USE test`, but does output the three `INSERT` statements following `USE db2`.

Row-based logging. `mysqlbinlog` outputs only entries that change tables belonging to *`db_name`*. The default database has no effect on this. Suppose that the binary log just described was created using row-based logging rather than statement-based logging. `mysqlbinlog --database=test` outputs only those entries that modify `t1` in the test database, regardless of whether `USE` was issued or what the default database is.

If a server is running with `binlog_format` set to `MIXED` and you want it to be possible to use `mysqlbinlog` with the `--database` option, you must ensure that tables that are modified are in the database selected by `USE`. (In particular, no cross-database updates should be used.)

When used together with the `--rewrite-db` option, the `--rewrite-db` option is applied first; then the `--database` option is applied, using the rewritten database name. The order in which the options are provided makes no difference in this regard.

- `--debug[=debug_options]`, `-# [debug_options]`

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,file_name`. The default is `d:t:o,/tmp/mysqlbinlog.trace`.

- `--debug-check`

  Print some debugging information when the program exits.

- `--debug-info`

  Print debugging information and memory and CPU usage statistics when the program exits.

- `--default-auth=plugin`

  A hint about the client-side authentication plugin to use. See Section 6.3.10, “Pluggable Authentication”.

- `--defaults-extra-file=file_name`

  Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *`file_name`* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

- `--defaults-file=file_name`

Use only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *file_name* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

Exception: Even with --defaults-file, client programs read .mylogin.cnf.

- --defaults-group-suffix=*str*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *str*. For example, mysqlbinlog normally reads the [client] and [mysqlbinlog] groups. If the --defaults-group-suffix=_other option is given, mysqlbinlog also reads the [client_other] and [mysqlbinlog_other] groups.

- --disable-log-bin, -D

  Disable binary logging. This is useful for avoiding an endless loop if you use the --to-last-log option and are sending the output to the same MySQL server. This option also is useful when restoring after a crash to avoid duplication of the statements you have logged.

  This option requires that you have the SYSTEM_VARIABLES_ADMIN or SUPER privilege. It causes mysqlbinlog to include a SET sql_log_bin = 0 statement in its output to disable binary logging of the remaining output. The SET statement is ineffective unless you have the SYSTEM_VARIABLES_ADMIN or SUPER privilege.

- --exclude-gtids=*gtid_set*

  Do not display any of the groups listed in the *gtid_set*.

- --force-if-open, -F

  Read binary log files even if they are open or were not closed properly.

- --force-read, -f

  With this option, if mysqlbinlog reads a binary log event that it does not recognize, it prints a warning, ignores the event, and continues. Without this option, mysqlbinlog stops if it reads such an event.

- --hexdump, -H

  Display a hex dump of the log in comments, as described in Section 4.6.8.1, “mysqlbinlog Hex Dump Format”. The hex output can be helpful for replication debugging.

- --host=*host_name*, -h *host_name*

  Get the binary log from the MySQL server on the given host.

- --idempotent

  Tell the MySQL Server to use idempotent mode while processing updates; this causes suppression of any duplicate-key or key-not-found errors that the server encounters in the current session while processing updates. This option may prove useful whenever it is desirable or necessary to replay one or more binary logs to a MySQL Server which may not contain all of the data to which the logs refer.

  The scope of effect for this option includes the current mysqlbinlog client and session only.

- --include-gtids=*gtid_set*

  Display only the groups listed in the *gtid_set*.

- --local-load=*dir_name*, -l *dir_name*

Prepare local temporary files for `LOAD DATA INFILE` in the specified directory.

Important

These temporary files are not automatically removed by `mysqlbinlog` or any other MySQL program.

- `--login-path=name`

  Read options from the named login path in the `.mylogin.cnf` login path file. A "login path" is an option group containing options that specify which MySQL server to connect to and which account to authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, "`mysql_config_editor` — MySQL Configuration Utility".

- `--no-defaults`

  Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read.

  The exception is that the `.mylogin.cnf` file, if it exists, is read in all cases. This permits passwords to be specified in a safer way than on the command line even when `--no-defaults` is used. (`.mylogin.cnf` is created by the `mysql_config_editor` utility. See Section 4.6.7, "`mysql_config_editor` — MySQL Configuration Utility".)

- `--offset=N`, `-o N`

  Skip the first `N` entries in the log.

- `--password[=password]`, `-p[password]`

  The password to use when connecting to the server. If you use the short option form (`-p`), you *cannot* have a space between the option and the password. If you omit the *`password`* value following the `--password` or `-p` option on the command line, `mysqlbinlog` prompts for one.

  Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, "End-User Guidelines for Password Security". You can use an option file to avoid giving the password on the command line.

- `--plugin-dir=dir_name`

  The directory in which to look for plugins. Specify this option if the `--default-auth` option is used to specify an authentication plugin but `mysqlbinlog` does not find it. See Section 6.3.10, "Pluggable Authentication".

- `--port=port_num`, `-P port_num`

  The TCP/IP port number to use for connecting to a remote server.

- `--print-defaults`

  Print the program name and all options that it gets from option files.

- `--print-table-metadata`

  Print table related metadata from the binary log. Configure the amount of table related metadata binary logged using `binlog-row-metadata`.

- `--protocol={TCP|SOCKET|PIPE|MEMORY}`

  The connection protocol to use for connecting to the server. It is useful when the other connection parameters normally would cause a protocol to be used other than the one you want. For details on the permissible values, see Section 4.2.2, "Connecting to the MySQL Server".

- `--raw`

  By default, `mysqlbinlog` reads binary log files and writes events in text format. The `--raw` option tells `mysqlbinlog` to write them in their original binary format. Its use requires that `--read-from-remote-server` also be used because the files are requested from a server. `mysqlbinlog` writes one output file for each file read from the server. The `--raw` option can be used to make a backup of a server's binary log. With the `--stop-never` option, the backup is "live" because `mysqlbinlog` stays connected to the server. By default, output files are written in the current directory with the same names as the original log files. Output file names can be modified using the `--result-file` option. For more information, see Section 4.6.8.3, "Using mysqlbinlog to Back Up Binary Log Files".

- `--read-from-remote-master=`*`type`*

  Read binary logs from a MySQL server with the `COM_BINLOG_DUMP` or `COM_BINLOG_DUMP_GTID` commands by setting the option value to either `BINLOG-DUMP-NON-GTIDS` or `BINLOG-DUMP-GTIDS`, respectively. If `--read-from-remote-master=BINLOG-DUMP-GTIDS` is combined with `--exclude-gtids`, transactions can be filtered out on the master, avoiding unnecessary network traffic.

  See also the description for `--read-from-remote-server`.

- `--read-from-remote-server`, `-R`

  Read the binary log from a MySQL server rather than reading a local log file. Any connection parameter options are ignored unless this option is given as well. These options are `--host`, `--password`, `--port`, `--protocol`, `--socket`, and `--user`.

  This option requires that the remote server be running. It works only for binary log files on the remote server, not relay log files.

  This option is like `--read-from-remote-master=BINLOG-DUMP-NON-GTIDS`.

- `--result-file=`*`name`*, `-r` *`name`*

  Without the `--raw` option, this option indicates the file to which `mysqlbinlog` writes text output. With `--raw`, `mysqlbinlog` writes one binary output file for each log file transferred from the server, writing them by default in the current directory using the same names as the original log file. In this case, the `--result-file` option value is treated as a prefix that modifies output file names.

- `--rewrite-db='`*`from_name->to_name`*`'`

  When reading from a row-based or statement-based log, rewrite all occurrences of *`from_name`* to *`to_name`*. Rewriting is done on the rows, for row-based logs, as well as on the `USE` clauses, for statement-based logs.

  Warning

  Statements in which table names are qualified with database names are not rewritten to use the new name when using this option.

  The rewrite rule employed as a value for this option is a string having the form `'`*`from_name->to_name`*`'`, as shown previously, and for this reason must be enclosed by quotation marks.

  To employ multiple rewrite rules, specify the option multiple times, as shown here:

```
shell> mysqlbinlog --rewrite-db='dbcurrent->dbold' --rewrite-db='dbtest->dbcurrent' \
                   binlog.00001 > /tmp/statements.sql
```

  When used together with the `--database` option, the `--rewrite-db` option is applied first; then `--database` option is applied, using the rewritten database name. The order in which the options are provided makes no difference in this regard.

This means that, for example, if `mysqlbinlog` is started with `--rewrite-db='mydb->yourdb' --database=yourdb`, then all updates to any tables in databases `mydb` and `yourdb` are included in the output. On the other hand, if it is started with `--rewrite-db='mydb->yourdb' --database=mydb`, then `mysqlbinlog` outputs no statements at all: since all updates to `mydb` are first rewritten as updates to `yourdb` before applying the `--database` option, there remain no updates that match `--database=mydb`.

- `--secure-auth`

  Do not send passwords to the server in old (pre-4.1) format. This prevents connections except for servers that use the newer password format.

  This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

  Note

  Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--server-id=`*`id`*

  Display only those events created by the server having the given server ID.

- `--set-charset=`*`charset_name`*

  Add a `SET NAMES` *`charset_name`* statement to the output to specify the character set to be used for processing log files.

- `--shared-memory-base-name=`*`name`*

  On Windows, the shared-memory name to use, for connections made using shared memory to a local server. The default value is `MYSQL`. The shared-memory name is case sensitive.

  The server must be started with the `--shared-memory` option to enable shared-memory connections.

- `--short-form`, `-s`

  Display only the statements contained in the log, without any extra information or row-based events. This is for testing only, and should not be used in production systems.

- `--skip-gtids[=(true|false)]`

  Do not display any GTIDs in the output. This is needed when writing to a dump file from one or more binary logs containing GTIDs, as shown in this example:

  ```
  shell> mysqlbinlog --skip-gtids binlog.000001 >  /tmp/dump.sql
  shell> mysqlbinlog --skip-gtids binlog.000002 >> /tmp/dump.sql
  shell> mysql -u root -p -e "source /tmp/dump.sql"
  ```

  The use of this option is otherwise not normally recommended in production.

- `--socket=`*`path`*, `-S` *`path`*

  For connections to `localhost`, the Unix socket file to use, or, on Windows, the name of the named pipe to use.

- `--ssl*`

Options that begin with `--ssl` specify whether to connect to the server using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--start-datetime=datetime`

  Start reading the binary log at the first event having a timestamp equal to or later than the *`datetime`* argument. The *`datetime`* value is relative to the local time zone on the machine where you run `mysqlbinlog`. The value should be in a format accepted for the `DATETIME` or `TIMESTAMP` data types. For example:

  ```
  shell> mysqlbinlog --start-datetime="2005-12-25 11:25:56" binlog.000003
  ```

  This option is useful for point-in-time recovery. See Section 7.3, “Example Backup and Recovery Strategy”.

- `--start-position=N`, `-j N`

  Start reading the binary log at the first event having a position equal to or greater than *`N`*. This option applies to the first log file named on the command line.

  This option is useful for point-in-time recovery. See Section 7.3, “Example Backup and Recovery Strategy”.

- `--stop-datetime=datetime`

  Stop reading the binary log at the first event having a timestamp equal to or later than the *`datetime`* argument. This option is useful for point-in-time recovery. See the description of the `--start-datetime` option for information about the *`datetime`* value.

  This option is useful for point-in-time recovery. See Section 7.3, “Example Backup and Recovery Strategy”.

- `--stop-never`

  This option is used with `--read-from-remote-server`. It tells `mysqlbinlog` to remain connected to the server. Otherwise `mysqlbinlog` exits when the last log file has been transferred from the server. `--stop-never` implies `--to-last-log`, so only the first log file to transfer need be named on the command line.

  `--stop-never` is commonly used with `--raw` to make a live binary log backup, but also can be used without `--raw` to maintain a continuous text display of log events as the server generates them.

- `--stop-never-slave-server-id=id`

  With `--stop-never`, `mysqlbinlog` reports a server ID of 65535 when it connects to the server. `--stop-never-slave-server-id` explicitly specifies the server ID to report. It can be used to avoid a conflict with the ID of a slave server or another `mysqlbinlog` process. See Section 4.6.8.4, “Specifying the mysqlbinlog Server ID”.

- `--stop-position=N`

  Stop reading the binary log at the first event having a position equal to or greater than *`N`*. This option applies to the last log file named on the command line.

  This option is useful for point-in-time recovery. See Section 7.3, “Example Backup and Recovery Strategy”.

- `--tls-version=protocol_list`

The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--to-last-log`, `-t`

  Do not stop at the end of the requested binary log from a MySQL server, but rather continue printing until the end of the last binary log. If you send the output to the same MySQL server, this may lead to an endless loop. This option requires `--read-from-remote-server`.

- `--user=user_name`, `-u user_name`

  The MySQL user name to use when connecting to a remote server.

- `--verbose`, `-v`

  Reconstruct row events and display them as commented SQL statements. If this option is given twice (by passing in either "-vv" or "--verbose --verbose"), the output includes comments to indicate column data types and some metadata, and row query log events if so configured.

  For examples that show the effect of `--base64-output` and `--verbose` on row event output, see Section 4.6.8.2, “mysqlbinlog Row Event Display”.

- `--verify-binlog-checksum`, `-c`

  Verify checksums in binary log files.

- `--version`, `-V`

  Display version information and exit.

  The `mysqlbinlog` version number shown when using this option is 3.4.

You can also set the following variable by using `--var_name=value` syntax:

- `open_files_limit`

  Specify the number of open file descriptors to reserve.

You can pipe the output of `mysqlbinlog` into the `mysql` client to execute the events contained in the binary log. This technique is used to recover from a crash when you have an old backup (see Section 7.5, “Point-in-Time (Incremental) Recovery Using the Binary Log”). For example:

```
shell> mysqlbinlog binlog.000001 | mysql -u root -p
```

Or:

```
shell> mysqlbinlog binlog.[0-9]* | mysql -u root -p
```

If the statements produced by `mysqlbinlog` may contain `BLOB` values, these may cause problems when `mysql` processes them. In this case, invoke `mysql` with the `--binary-mode` option.

You can also redirect the output of `mysqlbinlog` to a text file instead, if you need to modify the statement log first (for example, to remove statements that you do not want to execute for some reason). After editing the file, execute the statements that it contains by using it as input to the `mysql` program:

```
shell> mysqlbinlog binlog.000001 > tmpfile
shell> ... edit tmpfile ...
shell> mysql -u root -p < tmpfile
```

When mysqlbinlog is invoked with the --start-position option, it displays only those events with an offset in the binary log greater than or equal to a given position (the given position must match the start of one event). It also has options to stop and start when it sees an event with a given date and time. This enables you to perform point-in-time recovery using the --stop-datetime option (to be able to say, for example, “roll forward my databases to how they were today at 10:30 a.m.”).

If you have more than one binary log to execute on the MySQL server, the safe method is to process them all using a single connection to the server. Here is an example that demonstrates what may be *unsafe*:

```
shell> mysqlbinlog binlog.000001 | mysql -u root -p # DANGER!!
shell> mysqlbinlog binlog.000002 | mysql -u root -p # DANGER!!
```

Processing binary logs this way using multiple connections to the server causes problems if the first log file contains a CREATE TEMPORARY TABLE statement and the second log contains a statement that uses the temporary table. When the first mysql process terminates, the server drops the temporary table. When the second mysql process attempts to use the table, the server reports “unknown table.”

To avoid problems like this, use a *single* mysql process to execute the contents of all binary logs that you want to process. Here is one way to do so:

```
shell> mysqlbinlog binlog.000001 binlog.000002 | mysql -u root -p
```

Another approach is to write all the logs to a single file and then process the file:

```
shell> mysqlbinlog binlog.000001 >  /tmp/statements.sql
shell> mysqlbinlog binlog.000002 >> /tmp/statements.sql
shell> mysql -u root -p -e "source /tmp/statements.sql"
```

mysqlbinlog can produce output that reproduces a LOAD DATA INFILE operation without the original data file. mysqlbinlog copies the data to a temporary file and writes a LOAD DATA LOCAL INFILE statement that refers to the file. The default location of the directory where these files are written is system-specific. To specify a directory explicitly, use the --local-load option.

Because mysqlbinlog converts LOAD DATA INFILE statements to LOAD DATA LOCAL INFILE statements (that is, it adds LOCAL), both the client and the server that you use to process the statements must be configured with the LOCAL capability enabled. See Section 6.1.6, “Security Issues with LOAD DATA LOCAL”.

Warning

The temporary files created for LOAD DATA LOCAL statements are *not* automatically deleted because they are needed until you actually execute those statements. You should delete the temporary files yourself after you no longer need the statement log. The files can be found in the temporary file directory and have names like *original_file_name-#-#*.

4.6.8.1 mysqlbinlog Hex Dump Format

The --hexdump option causes mysqlbinlog to produce a hex dump of the binary log contents:

```
shell> mysqlbinlog --hexdump master-bin.000001
```

The hex output consists of comment lines beginning with #, so the output might look like this for the preceding command:

```
/*!40019 SET @@session.max_insert_delayed_threads=0*/;
/*!50003 SET @OLD_COMPLETION_TYPE=@@COMPLETION_TYPE,COMPLETION_TYPE=0*/;
# at 4
```

```
#051024 17:24:13 server id 1  end_log_pos 98
# Position  Timestamp   Type   Master ID        Size      Master Pos    Flags
# 00000004 9d fc 5c 43   0f   01 00 00 00   5e 00 00 00   62 00 00 00   00 00
# 00000017 04 00 35 2e 30 2e 31 35  2d 64 65 62 75 67 2d 6c |..5.0.15.debug.l|
# 00000027 6f 67 00 00 00 00 00 00  00 00 00 00 00 00 00 00 |og..............|
# 00000037 00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00 |................|
# 00000047 00 00 00 00 9d fc 5c 43  13 38 0d 00 08 00 12 00 |.......C.8......|
# 00000057 04 04 04 04 12 00 00 4b  00 04 1a               |.......K...|
#       Start: binlog v 4, server v 5.0.15-debug-log created 051024 17:24:13
#       at startup
ROLLBACK;
```

Hex dump output currently contains the elements in the following list. This format is subject to change. (For more information about binary log format, see MySQL Internals: The Binary Log.

- `Position`: The byte position within the log file.
- `Timestamp`: The event timestamp. In the example shown, `'9d fc 5c 43'` is the representation of `'051024 17:24:13'` in hexadecimal.
- `Type`: The event type code. In the example shown, `'0f'` indicates a `FORMAT_DESCRIPTION_EVENT`. The following table lists the possible type codes.

Type	Name	Meaning
00	UNKNOWN_EVENT	This event should never be present in the log.
01	START_EVENT_V3	This indicates the start of a log file written by MySQL 4 or earlier.
02	QUERY_EVENT	The most common type of events. These contain statements executed on the master.
03	STOP_EVENT	Indicates that master has stopped.
04	ROTATE_EVENT	Written when the master switches to a new log file.
05	INTVAR_EVENT	Used for AUTO_INCREMENT values or when the LAST_INSERT_ID() function is used in the statement.
06	LOAD_EVENT	Used for LOAD DATA INFILE in MySQL 3.23.
07	SLAVE_EVENT	Reserved for future use.
08	CREATE_FILE_EVENT	Used for LOAD DATA INFILE statements. This indicates the start of execution of such a statement. A temporary file is created on the slave. Used in MySQL 4 only.
09	APPEND_BLOCK_EVENT	Contains data for use in a LOAD DATA INFILE statement. The data is stored in the temporary file on the slave.
0a	EXEC_LOAD_EVENT	Used for LOAD DATA INFILE statements. The contents of the temporary file is stored in the table on the slave. Used in MySQL 4 only.
0b	DELETE_FILE_EVENT	Rollback of a LOAD DATA INFILE statement. The temporary file should be deleted on the slave.
0c	NEW_LOAD_EVENT	Used for LOAD DATA INFILE in MySQL 4 and earlier.
0d	RAND_EVENT	Used to send information about random values if the RAND() function is used in the statement.
0e	USER_VAR_EVENT	Used to replicate user variables.
0f	FORMAT_DESCRIPTION_EVENT	This indicates the start of a log file written by MySQL 5 or later.
10	XID_EVENT	Event indicating commit of an XA transaction.
11	BEGIN_LOAD_QUERY_EVENT	Used for LOAD DATA INFILE statements in MySQL 5 and later.

Type	Name	Meaning
12	EXECUTE_LOAD_QUERY_EVENT	Used for LOAD DATA INFILE statements in MySQL 5 and later.
13	TABLE_MAP_EVENT	Information about a table definition. Used in MySQL 5.1.5 and later.
14	PRE_GA_WRITE_ROWS_EVENT	Row data for a single table that should be created. Used in MySQL 5.1.5 to 5.1.17.
15	PRE_GA_UPDATE_ROWS_EVENT	Row data for a single table that needs to be updated. Used in MySQL 5.1.5 to 5.1.17.
16	PRE_GA_DELETE_ROWS_EVENT	Row data for a single table that should be deleted. Used in MySQL 5.1.5 to 5.1.17.
17	WRITE_ROWS_EVENT	Row data for a single table that should be created. Used in MySQL 5.1.18 and later.
18	UPDATE_ROWS_EVENT	Row data for a single table that needs to be updated. Used in MySQL 5.1.18 and later.
19	DELETE_ROWS_EVENT	Row data for a single table that should be deleted. Used in MySQL 5.1.18 and later.
1a	INCIDENT_EVENT	Something out of the ordinary happened. Added in MySQL 5.1.18.

- Master ID: The server ID of the master that created the event.
- Size: The size in bytes of the event.
- Master Pos: The position of the next event in the original master log file.
- Flags: 16 flags. The following flags are used. The others are reserved for future use.

Flag	Name	Meaning
01	LOG_EVENT_BINLOG_IN_USE_F	Log file correctly closed. (Used only in FORMAT_DESCRIPTION_EVENT.) If this flag is set (if the flags are, for example, '01 00') in a FORMAT_DESCRIPTION_EVENT, the log file has not been properly closed. Most probably this is because of a master crash (for example, due to power failure).
02		Reserved for future use.
04	LOG_EVENT_THREAD_SPECIFIC_	Set if the event is dependent on the connection it was executed in (for example, '04 00'), for example, if the event uses temporary tables.
08	LOG_EVENT_SUPPRESS_USE_F	Set in some circumstances when the event is not dependent on the default database.

4.6.8.2 mysqlbinlog Row Event Display

The following examples illustrate how mysqlbinlog displays row events that specify data modifications. These correspond to events with the WRITE_ROWS_EVENT, UPDATE_ROWS_EVENT, and DELETE_ROWS_EVENT type codes. The --base64-output=DECODE-ROWS and --verbose options may be used to affect row event output.

Suppose that the server is using row-based binary logging and that you execute the following sequence of statements:

```
CREATE TABLE t
(
```

```
  id   INT NOT NULL,
  name VARCHAR(20) NOT NULL,
  date DATE NULL
) ENGINE = InnoDB;

START TRANSACTION;
INSERT INTO t VALUES(1, 'apple', NULL);
UPDATE t SET name = 'pear', date = '2009-01-01' WHERE id = 1;
DELETE FROM t WHERE id = 1;
COMMIT;
```

By default, `mysqlbinlog` displays row events encoded as base-64 strings using `BINLOG` statements. Omitting extraneous lines, the output for the row events produced by the preceding statement sequence looks like this:

```
shell> mysqlbinlog log_file
...
# at 218
#080828 15:03:08 server id 1  end_log_pos 258   Write_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAANoAAAAAABEAAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBcBAAAAKAAAAAIBAAAQABEAAAAAAAEAA//8AQAAAAVhcHBsZQ==
'/*!*/;
...
# at 302
#080828 15:03:08 server id 1  end_log_pos 356   Update_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAC4BAAAAABEAAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBgBAAAANgAAAGQBAAAQABEAAAAAAAEAA////AEAAAAFYXBwbGX4AQAAAARwZWFyIbIP
'/*!*/;
...
# at 400
#080828 15:03:08 server id 1  end_log_pos 442   Delete_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAJABAAAAABEAAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBkBAAAAKgAAALoBAAAQABEAAAAAAAEAA//4AQAAAARwZWFyIbIP
'/*!*/;
```

To see the row events as comments in the form of “pseudo-SQL” statements, run `mysqlbinlog` with the `--verbose` or `-v` option. The output will contain lines beginning with `###`:

```
shell> mysqlbinlog -v log_file
...
# at 218
#080828 15:03:08 server id 1  end_log_pos 258   Write_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAANoAAAAAABEAAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBcBAAAAKAAAAAIBAAAQABEAAAAAAAEAA//8AQAAAAVhcHBsZQ==
'/*!*/;
### INSERT INTO test.t
### SET
###   @1=1
###   @2='apple'
###   @3=NULL
...
# at 302
#080828 15:03:08 server id 1  end_log_pos 356   Update_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAC4BAAAAABEAAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBgBAAAANgAAAGQBAAAQABEAAAAAAAEAA////AEAAAAFYXBwbGX4AQAAAARwZWFyIbIP
'/*!*/;
### UPDATE test.t
### WHERE
###   @1=1
```

```
###   @2='apple'
###   @3=NULL
### SET
###   @1=1
###   @2='pear'
###   @3='2009:01:01'
...
# at 400
#080828 15:03:08 server id 1  end_log_pos 442   Delete_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAJABAAAAABEAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBkBAAAAKgAAALoBAAAQABEAAAAAAAEAA//4AQAAAARwZWFyIbIP
'/*!*/;
### DELETE FROM test.t
### WHERE
###   @1=1
###   @2='pear'
###   @3='2009:01:01'
```

Specify `--verbose` or `-v` twice to also display data types and some metadata for each column. The output will contain an additional comment following each column change:

```
shell> mysqlbinlog -vv log_file
...
# at 218
#080828 15:03:08 server id 1  end_log_pos 258   Write_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAANoAAAAAABEAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBcBAAAAKAAAAAIBAAAQABEAAAAAAAEAA//8AQAAAAVhcHBsZQ==
'/*!*/;
### INSERT INTO test.t
### SET
###   @1=1 /* INT meta=0 nullable=0 is_null=0 */
###   @2='apple' /* VARSTRING(20) meta=20 nullable=0 is_null=0 */
###   @3=NULL /* VARSTRING(20) meta=0 nullable=1 is_null=1 */
...
# at 302
#080828 15:03:08 server id 1  end_log_pos 356   Update_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAC4BAAAAABEAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBgBAAAANgAAAGQBAAAQABEAAAAAAAEAA////AEAAAAFYXBwbGX4AQAAAARwZWFyIbIP
'/*!*/;
### UPDATE test.t
### WHERE
###   @1=1 /* INT meta=0 nullable=0 is_null=0 */
###   @2='apple' /* VARSTRING(20) meta=20 nullable=0 is_null=0 */
###   @3=NULL /* VARSTRING(20) meta=0 nullable=1 is_null=1 */
### SET
###   @1=1 /* INT meta=0 nullable=0 is_null=0 */
###   @2='pear' /* VARSTRING(20) meta=20 nullable=0 is_null=0 */
###   @3='2009:01:01' /* DATE meta=0 nullable=1 is_null=0 */
...
# at 400
#080828 15:03:08 server id 1  end_log_pos 442   Delete_rows: table id 17 flags: STMT_END_F

BINLOG '
fAS3SBMBAAAALAAAAJABAAAAABEAAAAAAAAABHRlc3QAAXQAAwMPCgIUAAQ=
fAS3SBkBAAAAKgAAALoBAAAQABEAAAAAAAEAA//4AQAAAARwZWFyIbIP
'/*!*/;
### DELETE FROM test.t
### WHERE
###   @1=1 /* INT meta=0 nullable=0 is_null=0 */
###   @2='pear' /* VARSTRING(20) meta=20 nullable=0 is_null=0 */
###   @3='2009:01:01' /* DATE meta=0 nullable=1 is_null=0 */
```

You can tell `mysqlbinlog` to suppress the `BINLOG` statements for row events by using the `--base64-output=DECODE-ROWS` option. This is similar to `--base64-output=NEVER` but does not

exit with an error if a row event is found. The combination of `--base64-output=DECODE-ROWS` and `--verbose` provides a convenient way to see row events only as SQL statements:

```
shell> mysqlbinlog -v --base64-output=DECODE-ROWS log_file
...
# at 218
#080828 15:03:08 server id 1  end_log_pos 258   Write_rows: table id 17 flags: STMT_END_F
### INSERT INTO test.t
### SET
###   @1=1
###   @2='apple'
###   @3=NULL
...
# at 302
#080828 15:03:08 server id 1  end_log_pos 356   Update_rows: table id 17 flags: STMT_END_F
### UPDATE test.t
### WHERE
###   @1=1
###   @2='apple'
###   @3=NULL
### SET
###   @1=1
###   @2='pear'
###   @3='2009:01:01'
...
# at 400
#080828 15:03:08 server id 1  end_log_pos 442   Delete_rows: table id 17 flags: STMT_END_F
### DELETE FROM test.t
### WHERE
###   @1=1
###   @2='pear'
###   @3='2009:01:01'
```

Note

You should not suppress `BINLOG` statements if you intend to re-execute `mysqlbinlog` output.

The SQL statements produced by `--verbose` for row events are much more readable than the corresponding `BINLOG` statements. However, they do not correspond exactly to the original SQL statements that generated the events. The following limitations apply:

- The original column names are lost and replaced by *`@N`*, where *`N`* is a column number.
- Character set information is not available in the binary log, which affects string column display:
  - There is no distinction made between corresponding binary and nonbinary string types (`BINARY` and `CHAR`, `VARBINARY` and `VARCHAR`, `BLOB` and `TEXT`). The output uses a data type of `STRING` for fixed-length strings and `VARSTRING` for variable-length strings.
  - For multibyte character sets, the maximum number of bytes per character is not present in the binary log, so the length for string types is displayed in bytes rather than in characters. For example, `STRING(4)` will be used as the data type for values from either of these column types:

    ```
    CHAR(4) CHARACTER SET latin1
    CHAR(2) CHARACTER SET ucs2
    ```

  - Due to the storage format for events of type `UPDATE_ROWS_EVENT`, `UPDATE` statements are displayed with the `WHERE` clause preceding the `SET` clause.

Proper interpretation of row events requires the information from the format description event at the beginning of the binary log. Because `mysqlbinlog` does not know in advance whether the rest of the log contains row events, by default it displays the format description event using a `BINLOG` statement in the initial part of the output.

If the binary log is known not to contain any events requiring a BINLOG statement (that is, no row events), the --base64-output=NEVER option can be used to prevent this header from being written.

4.6.8.3 Using mysqlbinlog to Back Up Binary Log Files

By default, mysqlbinlog reads binary log files and displays their contents in text format. This enables you to examine events within the files more easily and to re-execute them (for example, by using the output as input to mysql). mysqlbinlog can read log files directly from the local file system, or, with the --read-from-remote-server option, it can connect to a server and request binary log contents from that server. mysqlbinlog writes text output to its standard output, or to the file named as the value of the --result-file=*file_name* option if that option is given.

mysqlbinlog can read binary log files and write new files containing the same content—that is, in binary format rather than text format. This capability enables you to easily back up a binary log in its original format. mysqlbinlog can make a static backup, backing up a set of log files and stopping when the end of the last file is reached. It can also make a continuous ("live") backup, staying connected to the server when it reaches the end of the last log file and continuing to copy new events as they are generated. In continuous-backup operation, mysqlbinlog runs until the connection ends (for example, when the server exits) or mysqlbinlog is forcibly terminated. When the connection ends, mysqlbinlog does not wait and retry the connection, unlike a slave replication server. To continue a live backup after the server has been restarted, you must also restart mysqlbinlog.

Binary log backup requires that you invoke mysqlbinlog with two options at minimum:

- The --read-from-remote-server (or -R) option tells mysqlbinlog to connect to a server and request its binary log. (This is similar to a slave replication server connecting to its master server.)
- The --raw option tells mysqlbinlog to write raw (binary) output, not text output.

Along with --read-from-remote-server, it is common to specify other options: --host indicates where the server is running, and you may also need to specify connection options such as --user and --password.

Several other options are useful in conjunction with --raw:

- --stop-never: Stay connected to the server after reaching the end of the last log file and continue to read new events.
- --stop-never-slave-server-id=*id*: The server ID that mysqlbinlog reports to the server when --stop-never is used. The default is 65535. This can be used to avoid a conflict with the ID of a slave server or another mysqlbinlog process. See Section 4.6.8.4, "Specifying the mysqlbinlog Server ID".
- --result-file: A prefix for output file names, as described later.

To back up a server's binary log files with mysqlbinlog, you must specify file names that actually exist on the server. If you do not know the names, connect to the server and use the SHOW BINARY LOGS statement to see the current names. Suppose that the statement produces this output:

```
mysql> SHOW BINARY LOGS;
+---------------+-----------+
| Log_name      | File_size |
+---------------+-----------+
| binlog.000130 |     27459 |
| binlog.000131 |     13719 |
| binlog.000132 |     43268 |
+---------------+-----------+
```

With that information, you can use mysqlbinlog to back up the binary log to the current directory as follows (enter each command on a single line):

- To make a static backup of binlog.000130 through binlog.000132, use either of these commands:

```
mysqlbinlog --read-from-remote-server --host=host_name --raw
  binlog.000130 binlog.000131 binlog.000132

mysqlbinlog --read-from-remote-server --host=host_name --raw
  --to-last-log binlog.000130
```

The first command specifies every file name explicitly. The second names only the first file and uses `--to-last-log` to read through the last. A difference between these commands is that if the server happens to open `binlog.000133` before `mysqlbinlog` reaches the end of `binlog.000132`, the first command will not read it, but the second command will.

- To make a live backup in which `mysqlbinlog` starts with `binlog.000130` to copy existing log files, then stays connected to copy new events as the server generates them:

```
mysqlbinlog --read-from-remote-server --host=host_name --raw
  --stop-never binlog.000130
```

With `--stop-never`, it is not necessary to specify `--to-last-log` to read to the last log file because that option is implied.

Output File Naming

Without `--raw`, `mysqlbinlog` produces text output and the `--result-file` option, if given, specifies the name of the single file to which all output is written. With `--raw`, `mysqlbinlog` writes one binary output file for each log file transferred from the server. By default, `mysqlbinlog` writes the files in the current directory with the same names as the original log files. To modify the output file names, use the `--result-file` option. In conjunction with `--raw`, the `--result-file` option value is treated as a prefix that modifies the output file names.

Suppose that a server currently has binary log files named `binlog.000999` and up. If you use `mysqlbinlog --raw` to back up the files, the `--result-file` option produces output file names as shown in the following table. You can write the files to a specific directory by beginning the `--result-file` value with the directory path. If the `--result-file` value consists only of a directory name, the value must end with the pathname separator character. Output files are overwritten if they exist.

`--result-file` Option	**Output File Names**
`--result-file=x`	`xbinlog.000999` and up
`--result-file=/tmp/`	`/tmp/binlog.000999` and up
`--result-file=/tmp/x`	`/tmp/xbinlog.000999` and up

Example: mysqldump + mysqlbinlog for Backup and Restore

The following example describes a simple scenario that shows how to use `mysqldump` and `mysqlbinlog` together to back up a server's data and binary log, and how to use the backup to restore the server if data loss occurs. The example assumes that the server is running on host *`host_name`* and its first binary log file is named `binlog.000999`. Enter each command on a single line.

Use `mysqlbinlog` to make a continuous backup of the binary log:

```
mysqlbinlog --read-from-remote-server --host=host_name --raw
  --stop-never binlog.000999
```

Use `mysqldump` to create a dump file as a snapshot of the server's data. Use `--all-databases`, `--events`, and `--routines` to back up all data, and `--master-data=2` to include the current binary log coordinates in the dump file.

```
mysqldump --host=host_name --all-databases --events --routines --master-data=2> dump_file
```

Execute the mysqldump command periodically to create newer snapshots as desired.

If data loss occurs (for example, if the server crashes), use the most recent dump file to restore the data:

```
mysql --host=host_name -u root -p < dump_file
```

Then use the binary log backup to re-execute events that were written after the coordinates listed in the dump file. Suppose that the coordinates in the file look like this:

```
-- CHANGE MASTER TO MASTER_LOG_FILE='binlog.001002', MASTER_LOG_POS=27284;
```

If the most recent backed-up log file is named binlog.001004, re-execute the log events like this:

```
mysqlbinlog --start-position=27284 binlog.001002 binlog.001003 binlog.001004
  | mysql --host=host_name -u root -p
```

You might find it easier to copy the backup files (dump file and binary log files) to the server host to make it easier to perform the restore operation, or if MySQL does not allow remote root access.

4.6.8.4 Specifying the mysqlbinlog Server ID

When invoked with the --read-from-remote-server option, mysqlbinlog connects to a MySQL server, specifies a server ID to identify itself, and requests binary log files from the server. You can use mysqlbinlog to request log files from a server in several ways:

- Specify an explicitly named set of files: For each file, mysqlbinlog connects and issues a Binlog dump command. The server sends the file and disconnects. There is one connection per file.
- Specify the beginning file and --to-last-log: mysqlbinlog connects and issues a Binlog dump command for all files. The server sends all files and disconnects.
- Specify the beginning file and --stop-never (which implies --to-last-log): mysqlbinlog connects and issues a Binlog dump command for all files. The server sends all files, but does not disconnect after sending the last one.

With --read-from-remote-server only, mysqlbinlog connects using a server ID of 0, which tells the server to disconnect after sending the last requested log file.

With --read-from-remote-server and --stop-never, mysqlbinlog connects using a nonzero server ID, so the server does not disconnect after sending the last log file. The server ID is 65535 by default, but this can be changed with --stop-never-slave-server-id.

Thus, for the first two ways of requesting files, the server disconnects because mysqlbinlog specifies a server ID of 0. It does not disconnect if --stop-never is given because mysqlbinlog specifies a nonzero server ID.

4.6.9 mysqldumpslow — Summarize Slow Query Log Files

The MySQL slow query log contains information about queries that take a long time to execute (see Section 5.4.5, “The Slow Query Log”). mysqldumpslow parses MySQL slow query log files and prints a summary of their contents.

Normally, mysqldumpslow groups queries that are similar except for the particular values of number and string data values. It “abstracts” these values to N and 'S' when displaying summary output. The -a and -n options can be used to modify value abstracting behavior.

Invoke mysqldumpslow like this:

```
shell> mysqldumpslow [options] [log_file ...]
```

mysqldumpslow supports the following options.

Table 4.19 mysqldumpslow Options

Format	Description
-a	Do not abstract all numbers to N and strings to S
-n	Abstract numbers with at least the specified digits
--debug	Write debugging information
-g	Only consider statements that match the pattern
--help	Display help message and exit
-h	Host name of the server in the log file name
-i	Name of the server instance
-l	Do not subtract lock time from total time
-r	Reverse the sort order
-s	How to sort output
-t	Display only first num queries
--verbose	Verbose mode

- `--help`

  Display a help message and exit.

- `-a`

  Do not abstract all numbers to `N` and strings to `'S'`.

- `--debug`, `-d`

  Run in debug mode.

- `-g pattern`

  Consider only queries that match the (`grep`-style) pattern.

- `-h host_name`

  Host name of MySQL server for `*-slow.log` file name. The value can contain a wildcard. The default is `*` (match all).

- `-i name`

  Name of server instance (if using `mysql.server` startup script).

- `-l`

  Do not subtract lock time from total time.

- `-n N`

  Abstract numbers with at least `N` digits within names.

- `-r`

  Reverse the sort order.

- `-s sort_type`

  How to sort the output. The value of `sort_type` should be chosen from the following list:

- `t`, `at`: Sort by query time or average query time
- `l`, `al`: Sort by lock time or average lock time
- `r`, `ar`: Sort by rows sent or average rows sent
- `c`: Sort by count

By default, `mysqldumpslow` sorts by average query time (equivalent to `-s at`).

- `-t N`

  Display only the first `N` queries in the output.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does.

Example of usage:

```
shell> mysqldumpslow

Reading mysql slow query log from /usr/local/mysql/data/mysqld51-apple-slow.log
Count: 1  Time=4.32s (4s)  Lock=0.00s (0s)  Rows=0.0 (0), root[root]@localhost
 insert into t2 select * from t1

Count: 3  Time=2.53s (7s)  Lock=0.00s (0s)  Rows=0.0 (0), root[root]@localhost
 insert into t2 select * from t1 limit N

Count: 3  Time=2.13s (6s)  Lock=0.00s (0s)  Rows=0.0 (0), root[root]@localhost
 insert into t1 select * from t1
```

4.7 MySQL Program Development Utilities

This section describes some utilities that you may find useful when developing MySQL programs.

In shell scripts, you can use the `my_print_defaults` program to parse option files and see what options would be used by a given program. The following example shows the output that `my_print_defaults` might produce when asked to show the options found in the `[client]` and `[mysql]` groups:

```
shell> my_print_defaults client mysql
--port=3306
--socket=/tmp/mysql.sock
--no-auto-rehash
```

Note for developers: Option file handling is implemented in the C client library simply by processing all options in the appropriate group or groups before any command-line arguments. This works well for programs that use the last instance of an option that is specified multiple times. If you have a C or C++ program that handles multiply specified options this way but that doesn't read option files, you need add only two lines to give it that capability. Check the source code of any of the standard MySQL clients to see how to do this.

Several other language interfaces to MySQL are based on the C client library, and some of them provide a way to access option file contents. These include Perl and Python. For details, see the documentation for your preferred interface.

4.7.1 `mysql_config` — Display Options for Compiling Clients

`mysql_config` provides you with useful information for compiling your MySQL client and connecting it to MySQL. It is a shell script, so it is available only on Unix and Unix-like systems.

Note

pkg-config can be used as an alternative to mysql_config for obtaining information such as compiler flags or link libraries required to compile MySQL applications. For more information, see Section 27.7.4.2, “Building C API Client Programs Using pkg-config”.

mysql_config supports the following options.

- --cflags

  C Compiler flags to find include files and critical compiler flags and defines used when compiling the libmysqlclient library. The options returned are tied to the specific compiler that was used when the library was created and might clash with the settings for your own compiler. Use --include for more portable options that contain only include paths.

- --cxxflags

  Like --cflags, but for C++ compiler flags.

- --include

  Compiler options to find MySQL include files.

- --libs

  Libraries and options required to link with the MySQL client library.

- --libs_r

  Libraries and options required to link with the thread-safe MySQL client library. In MySQL 8.0, all client libraries are thread-safe, so this option need not be used. The --libs option can be used in all cases.

- --plugindir

  The default plugin directory path name, defined when configuring MySQL.

- --port

  The default TCP/IP port number, defined when configuring MySQL.

- --socket

  The default Unix socket file, defined when configuring MySQL.

- --variable=*var_name*

  Display the value of the named configuration variable. Permitted *var_name* values are pkgincludedir (the header file directory), pkglibdir (the library directory), and plugindir (the plugin directory).

- --version

  Version number for the MySQL distribution.

If you invoke mysql_config with no options, it displays a list of all options that it supports, and their values:

```
shell> mysql_config
Usage: /usr/local/mysql/bin/mysql_config [options]
Options:
  --cflags         [-I/usr/local/mysql/include/mysql -mcpu=pentiumpro]
```

```
  --cxxflags       [-I/usr/local/mysql/include/mysql -mcpu=pentiumpro]
  --include        [-I/usr/local/mysql/include/mysql]
  --libs           [-L/usr/local/mysql/lib/mysql -lmysqlclient
                    -lpthread -lm -lrt -lssl -lcrypto -ldl]
  --libs_r         [-L/usr/local/mysql/lib/mysql -lmysqlclient_r
                    -lpthread -lm -lrt -lssl -lcrypto -ldl]
  --plugindir      [/usr/local/mysql/lib/plugin]
  --socket         [/tmp/mysql.sock]
  --port           [3306]
  --version        [5.8.0-m17]
  --variable=VAR   VAR is one of:
          pkgincludedir [/usr/local/mysql/include]
          pkglibdir     [/usr/local/mysql/lib]
          plugindir     [/usr/local/mysql/lib/plugin]
```

You can use `mysql_config` within a command line using backticks to include the output that it produces for particular options. For example, to compile and link a MySQL client program, use `mysql_config` as follows:

```
gcc -c `mysql_config --cflags` progname.c
gcc -o progname progname.o `mysql_config --libs`
```

4.7.2 `my_print_defaults` — Display Options from Option Files

`my_print_defaults` displays the options that are present in option groups of option files. The output indicates what options will be used by programs that read the specified option groups. For example, the `mysqlcheck` program reads the `[mysqlcheck]` and `[client]` option groups. To see what options are present in those groups in the standard option files, invoke `my_print_defaults` like this:

```
shell> my_print_defaults mysqlcheck client
--user=myusername
--password=secret
--host=localhost
```

The output consists of options, one per line, in the form that they would be specified on the command line.

`my_print_defaults` supports the following options.

- `--help`, `-?`

  Display a help message and exit.

- `--config-file=`*`file_name`*, `--defaults-file=`*`file_name`*, `-c` *`file_name`*

  Read only the given option file.

- `--debug=`*`debug_options`*, `-#` *`debug_options`*

  Write a debugging log. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:o,/tmp/my_print_defaults.trace`.

- `--defaults-extra-file=`*`file_name`*, `--extra-file=`*`file_name`*, `-e` *`file_name`*

  Read this option file after the global option file but (on Unix) before the user option file.

- `--defaults-group-suffix=`*`suffix`*, `-g` *`suffix`*

  In addition to the groups named on the command line, read groups that have the given suffix.

- `--login-path=`*`name`*, `-l` *`name`*

  Read options from the named login path in the `.mylogin.cnf` login path file. A “login path” is an option group containing options that specify which MySQL server to connect to and which account to

authenticate as. To create or modify a login path file, use the `mysql_config_editor` utility. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- `--no-defaults`, `-n`

  Return an empty string.

- `--show`, `-s`

  `my_print_defaults` masks passwords by default. Use this option to display passwords in cleartext.

- `--verbose`, `-v`

  Verbose mode. Print more information about what the program does.

- `--version`, `-V`

  Display version information and exit.

4.7.3 `resolve_stack_dump` — Resolve Numeric Stack Trace Dump to Symbols

`resolve_stack_dump` resolves a numeric stack dump to symbols.

Invoke `resolve_stack_dump` like this:

```
shell> resolve_stack_dump [options] symbols_file [numeric_dump_file]
```

The symbols file should include the output from the `nm --numeric-sort mysqld` command. The numeric dump file should contain a numeric stack track from `mysqld`. If no numeric dump file is named on the command line, the stack trace is read from the standard input.

`resolve_stack_dump` supports the following options.

- `--help`, `-h`

  Display a help message and exit.

- `--numeric-dump-file=`*`file_name`*, `-n` *`file_name`*

  Read the stack trace from the given file.

- `--symbols-file=`*`file_name`*, `-s` *`file_name`*

  Use the given symbols file.

- `--version`, `-V`

  Display version information and exit.

For more information, see Section 28.5.1.5, “Using a Stack Trace”.

4.8 Miscellaneous Programs

4.8.1 `lz4_decompress` — Decompress mysqlpump LZ4-Compressed Output

The `lz4_decompress` utility decompresses `mysqlpump` output that was created using LZ4 compression.

Note

If MySQL was configured with the `-DWITH_LZ4=system` option, `lz4_decompress` is not built. In this case, the system `lz4` command can be used instead.

Invoke `lz4_decompress` like this:

```
shell> lz4_decompress input_file output_file
```

Example:

```
shell> mysqlpump --compress-output=LZ4 > dump.lz4
shell> lz4_decompress dump.lz4 dump.txt
```

To see a help message, invoke `lz4_decompress` with no arguments.

To decompress `mysqlpump` ZLIB-compressed output, use `zlib_decompress`. See Section 4.8.4, “`zlib_decompress` — Decompress mysqlpump ZLIB-Compressed Output”.

4.8.2 `perror` — Explain Error Codes

For most system errors, MySQL displays, in addition to an internal text message, the system error code in one of the following styles:

```
message ... (errno: #)
message ... (Errcode: #)
```

You can find out what the error code means by examining the documentation for your system or by using the `perror` utility.

`perror` prints a description for a system error code or for a storage engine (table handler) error code.

Invoke `perror` like this:

```
shell> perror [options] errorcode ...
```

Example:

```
shell> perror 13 64
OS error code  13:  Permission denied
OS error code  64:  Machine is not on the network
```

To obtain the error message for a MySQL Cluster error code, invoke `perror` with the `--ndb` option:

```
shell> perror --ndb errorcode
```

The meaning of system error messages may be dependent on your operating system. A given error code may mean different things on different operating systems.

`perror` supports the following options.

- `--help`, `--info`, `-I`, `-?`

  Display a help message and exit.

- `--ndb`

  Print the error message for a MySQL Cluster error code.

- `--silent`, `-s`

  Silent mode. Print only the error message.

- `--verbose`, `-v`

  Verbose mode. Print error code and message. This is the default behavior.

- `--version`, `-V`

  Display version information and exit.

4.8.3 `resolveip` — Resolve Host name to IP Address or Vice Versa

The `resolveip` utility resolves host names to IP addresses and vice versa.

Invoke `resolveip` like this:

```
shell> resolveip [options] {host_name|ip-addr} ...
```

`resolveip` supports the following options.

- `--help`, `--info`, `-?`, `-I`

  Display a help message and exit.

- `--silent`, `-s`

  Silent mode. Produce less output.

- `--version`, `-V`

  Display version information and exit.

4.8.4 `zlib_decompress` — Decompress mysqlpump ZLIB-Compressed Output

The `zlib_decompress` utility decompresses `mysqlpump` output that was created using ZLIB compression.

Note

If MySQL was configured with the `-DWITH_ZLIB=system` option, `zlib_decompress` is not built. In this case, the system `openssl zlib` command can be used instead.

Invoke `zlib_decompress` like this:

```
shell> zlib_decompress input_file output_file
```

Example:

```
shell> mysqlpump --compress-output=ZLIB > dump.zlib
shell> zlib_decompress dump.zlib dump.txt
```

To see a help message, invoke `zlib_decompress` with no arguments.

To decompress `mysqlpump` LZ4-compressed output, use `lz4_decompress`. See Section 4.8.1, “`lz4_decompress` — Decompress mysqlpump LZ4-Compressed Output”.

4.9 MySQL Program Environment Variables

This section lists environment variables that are used directly or indirectly by MySQL. Most of these can also be found in other places in this manual.

Options on the command line take precedence over values specified in option files and environment variables, and values in option files take precedence over values in environment variables. In many cases, it is preferable to use an option file instead of environment variables to modify the behavior of MySQL. See Section 4.2.6, “Using Option Files”.

Variable	Description
`CXX`	The name of your C++ compiler (for running `CMake`).
`CC`	The name of your C compiler (for running `CMake`).
`DBI_USER`	The default user name for Perl DBI.
`DBI_TRACE`	Trace options for Perl DBI.
`HOME`	The default path for the `mysql` history file is `$HOME/.mysql_history`.
`LD_RUN_PATH`	Used to specify the location of `libmysqlclient.so`.
`LIBMYSQL_ENABLE_CLEARTEXT_PLUGIN`	Enable `mysql_clear_password` authentication plugin; see Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.
`LIBMYSQL_PLUGIN_DIR`	Directory in which to look for client plugins.
`LIBMYSQL_PLUGINS`	Client plugins to preload.
`MYSQL_DEBUG`	Debug trace options when debugging.
`MYSQL_GROUP_SUFFIX`	Option group suffix value (like specifying `--defaults-group-suffix`).
`MYSQL_HISTFILE`	The path to the `mysql` history file. If this variable is set, its value overrides the default for `$HOME/.mysql_history`.
`MYSQL_HISTIGNORE`	Patterns specifying statements that `mysql` should not log to `$HOME/.mysql_history`, or `syslog` if `--syslog` is given.
`MYSQL_HOME`	The path to the directory in which the server-specific `my.cnf` file resides.
`MYSQL_HOST`	The default host name used by the `mysql` command-line client.
`MYSQL_PS1`	The command prompt to use in the `mysql` command-line client.
`MYSQL_PWD`	The default password when connecting to `mysqld`. Using this is insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”.
`MYSQL_TCP_PORT`	The default TCP/IP port number.
`MYSQL_TEST_LOGIN_FILE`	The name of the `.mylogin.cnf` login path file.
`MYSQL_TEST_TRACE_CRASH`	Whether the test protocol trace plugin crashes clients. See note following table.
`MYSQL_TEST_TRACE_DEBUG`	Whether the test protocol trace plugin produces output. See note following table.
`MYSQL_UNIX_PORT`	The default Unix socket file name; used for connections to `localhost`.
`PATH`	Used by the shell to find MySQL programs.

Variable	Description
PKG_CONFIG_PATH	Location of mysqlclient.pc pkg-config file. See note following table.
TMPDIR	The directory in which temporary files are created.
TZ	This should be set to your local time zone. See Section B.5.3.7, “Time Zone Problems”.
UMASK	The user-file creation mode when creating files. See note following table.
UMASK_DIR	The user-directory creation mode when creating directories. See note following table.
USER	The default user name on Windows when connecting to mysqld.

For information about the mysql history file, see Section 4.5.1.3, “mysql Logging”.

MYSQL_TEST_LOGIN_FILE is the path name of the login path file (the file created by mysql_config_editor). If not set, the default value is %APPDATA%\MySQL\.mylogin.cnf directory on Windows and $HOME/.mylogin.cnf on non-Windows systems. See Section 4.6.7, “mysql_config_editor — MySQL Configuration Utility”.

The MYSQL_TEST_TRACE_DEBUG and MYSQL_TEST_TRACE_CRASH variables control the test protocol trace client plugin, if MySQL is built with that plugin enabled. For more information, see Using the Test Protocol Trace Plugin.

The default UMASK and UMASK_DIR values are 0640 and 0750, respectively. MySQL assumes that the value for UMASK or UMASK_DIR is in octal if it starts with a zero. For example, setting UMASK=0600 is equivalent to UMASK=384 because 0600 octal is 384 decimal.

The UMASK and UMASK_DIR variables, despite their names, are used as modes, not masks:

- If UMASK is set, mysqld uses ($UMASK | 0600) as the mode for file creation, so that newly created files have a mode in the range from 0600 to 0666 (all values octal).
- If UMASK_DIR is set, mysqld uses ($UMASK_DIR | 0700) as the base mode for directory creation, which then is AND-ed with ~(~$UMASK & 0666), so that newly created directories have a mode in the range from 0700 to 0777 (all values octal). The AND operation may remove read and write permissions from the directory mode, but not execute permissions.

It may be necessary to set PKG_CONFIG_PATH if you use pkg-config for building MySQL programs. See Section 27.7.4.2, “Building C API Client Programs Using pkg-config”.

Chapter 5 MySQL Server Administration

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MySQL Server (`mysqld`) is the main program that does most of the work in a MySQL installation. This chapter provides an overview of MySQL Server and covers general server administration:

- Server configuration
- The data directory, particularly the `mysql` system database
- The server log files
- Management of multiple servers on a single machine

For additional information on administrative topics, see also:

- Chapter 6, *Security*
- Chapter 7, *Backup and Recovery*
- Chapter 18, *Replication*

5.1 The MySQL Server

`mysqld` is the MySQL server. The following discussion covers these MySQL server configuration topics:

- Startup options that the server supports. You can specify these options on the command line, through configuration files, or both.
- Server system variables. These variables reflect the current state and values of the startup options, some of which can be modified while the server is running.
- Server status variables. These variables contain counters and statistics about runtime operation.
- How to set the server SQL mode. This setting modifies certain aspects of SQL syntax and semantics, for example for compatibility with code from other database systems, or to control the error handling for particular situations.
- The server shutdown process. There are performance and reliability considerations depending on the type of table (transactional or nontransactional) and whether you use replication.

For listings of MySQL server variables and options that have been added, deprecated, or removed in MySQL 8.0, see Section 1.5, “Server and Status Variables and Options Added, Deprecated, or Removed in MySQL 8.0”.

Note

Not all storage engines are supported by all MySQL server binaries and configurations. To find out how to determine which storage engines your MySQL server installation supports, see Section 13.7.5.16, “SHOW ENGINES Syntax”.

5.1.1 Configuring the Server

The MySQL server, `mysqld`, has many command options and system variables that can be set at startup to configure its operation. To determine the command option and system variable values used by the server, execute this command:

```
shell> mysqld --verbose --help
```

The command produces a list of all `mysqld` options and configurable system variables. Its output includes the default option and variable values and looks something like this:

```
abort-slave-event-count           0
allow-suspicious-udfs             FALSE
archive                           ON
auto-increment-increment          1
auto-increment-offset             1
autocommit                        TRUE
automatic-sp-privileges           TRUE
avoid-temporal-upgrade            FALSE
back-log                          80
basedir                           /home/jon/bin/mysql-8.0/
...
tmpdir                            /tmp
transaction-alloc-block-size      8192
transaction-isolation             REPEATABLE-READ
transaction-prealloc-size         4096
transaction-read-only             FALSE
transaction-write-set-extraction  OFF
updatable-views-with-limit        YES
validate-user-plugins             TRUE
verbose                           TRUE
wait-timeout                      28800
```

To see the current system variable values used by the server as it runs, connect to it and execute this statement:

```
mysql> SHOW VARIABLES;
```

To see some statistical and status indicators for a running server, execute this statement:

```
mysql> SHOW STATUS;
```

System variable and status information also is available using the `mysqladmin` command:

```
shell> mysqladmin variables
shell> mysqladmin extended-status
```

For a full description of all command options, system variables, and status variables, see these sections:

- Section 5.1.4, “Server Command Options”
- Section 5.1.5, “Server System Variables”
- Section 5.1.7, “Server Status Variables”

More detailed monitoring information is available from the Performance Schema; see Chapter 25, *MySQL Performance Schema*.

If you specify an option on the command line for `mysqld` or `mysqld_safe`, it remains in effect only for that invocation of the server. To use the option every time the server runs, put it in an option file. See Section 4.2.6, “Using Option Files”.

5.1.2 Server Configuration Defaults

The MySQL server has many operating parameters, which you can change at server startup using command-line options or configuration files (option files). It is also possible to change many parameters at runtime. For general instructions on setting parameters at startup or runtime, see Section 5.1.4, “Server Command Options”, and Section 5.1.5, “Server System Variables”.

On Windows, MySQL Installer interacts with the user and creates a file named `my.ini` in the base installation directory as the default option file.

Note

On Windows, the `.ini` or `.cnf` option file extension might not be displayed.

On any platform, after completing the installation process, you can edit the default option file at any time to modify the parameters used by the server. For example, to use a parameter setting in the file that is commented with a `#` character at the beginning of the line, remove the `#`, and modify the parameter value if necessary. To disable a setting, either add a `#` to the beginning of the line or remove it.

For additional information about option file format and syntax, see Section 4.2.6, “Using Option Files”.

5.1.3 Server Option and Variable Reference

The following table provides a list of all the command line options, server and status variables applicable within `mysqld`.

The table lists command-line options (Cmd-line), options valid in configuration files (Option file), server system variables (System Var), and status variables (Status var) in one unified list, with notification of where each option/variable is valid. If a server option set on the command line or in an option file differs from the name of the corresponding server system or status variable, the variable name is noted

immediately below the corresponding option. For status variables, the scope of the variable is shown (Scope) as either global, session, or both. Please see the corresponding sections for details on setting and using the options and variables. Where appropriate, a direct link to further information on the item as available.

Table 5.1 Option/Variable Summary

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
abort-slave-event-count	Yes	Yes				
Aborted_clients				Yes	Global	No
Aborted_connects				Yes	Global	No
Acl_cache_items_count				Yes	Global	No
allow-suspicious-udfs	Yes	Yes				
ansi	Yes	Yes				
auto_generate_certs	Yes	Yes	Yes		Global	No
auto_increment_increment			Yes		Both	Yes
auto_increment_offset			Yes		Both	Yes
autocommit	Yes	Yes	Yes		Both	Yes
automatic_sp_privileges			Yes		Global	Yes
avoid_temporal_upgrade	Yes	Yes	Yes		Global	Yes
back_log			Yes		Global	No
basedir	Yes	Yes	Yes		Global	No
big-tables	Yes	Yes			Both	Yes
- *Variable*: big_tables			Yes		Both	Yes
bind-address	Yes	Yes			Global	No
- *Variable*: bind_address			Yes		Global	No
Binlog_cache_disk_use				Yes	Global	No
binlog_cache_size	Yes	Yes	Yes		Global	Yes
Binlog_cache_use				Yes	Global	No
binlog-checksum	Yes	Yes				
binlog_checksum			Yes		Global	Yes
binlog_direct_non_transactional_updates	Yes		Yes		Both	Yes
binlog-do-db	Yes	Yes				
binlog_error_action	Yes	Yes	Yes		Both	Yes
binlog_expire_logs_seconds	Yes	Yes	Yes		Global	Yes
binlog-format	Yes	Yes			Both	Yes
- *Variable*: binlog_format			Yes		Both	Yes
binlog_group_commit_sync_delay	Yes	Yes	Yes		Global	Yes
binlog_group_commit_sync_no_delay_count	Yes		Yes		Global	Yes
binlog_gtid_simple_recovery	Yes	Yes	Yes		Global	No
binlog-ignore-db	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
binlog_max_flush_queue_time			Yes		Global	Yes
binlog_order_commits			Yes		Global	Yes
binlog-row-event-max-size	Yes	Yes				
binlog_row_image	Yes	Yes	Yes		Both	Yes
binlog_row_metadata	Yes	Yes	Yes		Global	Yes
binlog-rows-query-log-events	Yes	Yes				
- *Variable*: binlog_rows_query_log_events						
binlog_rows_query_log_events			Yes		Both	Yes
Binlog_stmt_cache_disk_use				Yes	Global	No
binlog_stmt_cache_size	Yes	Yes	Yes		Global	Yes
Binlog_stmt_cache_use				Yes	Global	No
binlog_transaction_dependency_history_size	Yes		Yes		Global	Yes
binlog_transaction_dependency_tracking	Yes		Yes		Global	Yes
binlogging_impossible_mode	Yes	Yes	Yes		Both	Yes
block_encryption_mode	Yes	Yes	Yes		Both	Yes
bulk_insert_buffer_size	Yes	Yes	Yes		Both	Yes
Bytes_received				Yes	Both	No
Bytes_sent				Yes	Both	No
character_set_client			Yes		Both	Yes
character-set-client-handshake	Yes	Yes				
character_set_connection			Yes		Both	Yes
character_set_database[a]			Yes		Both	Yes
character-set-filesystem	Yes	Yes			Both	Yes
- *Variable*: character_set_filesystem			Yes		Both	Yes
character_set_results			Yes		Both	Yes
character-set-server	Yes	Yes			Both	Yes
- *Variable*: character_set_server			Yes		Both	Yes
character_set_system			Yes		Global	No
character-sets-dir	Yes	Yes			Global	No
- *Variable*: character_sets_dir			Yes		Global	No
check_proxy_users	Yes	Yes	Yes		Global	Yes
chroot	Yes	Yes				
collation_connection			Yes		Both	Yes
collation_database[b]			Yes		Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
collation-server	Yes	Yes			Both	Yes
- *Variable*: collation_server			Yes		Both	Yes
Com_admin_commands				Yes	Both	No
Com_alter_db				Yes	Both	No
Com_alter_event				Yes	Both	No
Com_alter_function				Yes	Both	No
Com_alter_procedure				Yes	Both	No
Com_alter_server				Yes	Both	No
Com_alter_table				Yes	Both	No
Com_alter_tablespace				Yes	Both	No
Com_alter_user				Yes	Both	No
Com_alter_user_default_role				Yes	Global	No
Com_analyze				Yes	Both	No
Com_assign_to_keycache				Yes	Both	No
Com_begin				Yes	Both	No
Com_binlog				Yes	Both	No
Com_call_procedure				Yes	Both	No
Com_change_db				Yes	Both	No
Com_change_master				Yes	Both	No
Com_change_repl_filter				Yes	Both	No
Com_check				Yes	Both	No
Com_checksum				Yes	Both	No
Com_commit				Yes	Both	No
Com_create_db				Yes	Both	No
Com_create_event				Yes	Both	No
Com_create_function				Yes	Both	No
Com_create_index				Yes	Both	No
Com_create_procedure				Yes	Both	No
Com_create_role				Yes	Global	No
Com_create_server				Yes	Both	No
Com_create_table				Yes	Both	No
Com_create_trigger				Yes	Both	No
Com_create_udf				Yes	Both	No
Com_create_user				Yes	Both	No
Com_create_view				Yes	Both	No
Com_dealloc_sql				Yes	Both	No
Com_delete				Yes	Both	No
Com_delete_multi				Yes	Both	No
Com_do				Yes	Both	No
Com_drop_db				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Com_drop_event				Yes	Both	No
Com_drop_function				Yes	Both	No
Com_drop_index				Yes	Both	No
Com_drop_procedure				Yes	Both	No
Com_drop_role				Yes	Global	No
Com_drop_server				Yes	Both	No
Com_drop_table				Yes	Both	No
Com_drop_trigger				Yes	Both	No
Com_drop_user				Yes	Both	No
Com_drop_view				Yes	Both	No
Com_empty_query				Yes	Both	No
Com_execute_sql				Yes	Both	No
Com_explain_other				Yes	Both	No
Com_flush				Yes	Both	No
Com_get_diagnostics				Yes	Both	No
Com_grant				Yes	Both	No
Com_grant_roles				Yes	Global	No
Com_group_replication_start				Yes	Global	No
Com_group_replication_stop				Yes	Global	No
Com_ha_close				Yes	Both	No
Com_ha_open				Yes	Both	No
Com_ha_read				Yes	Both	No
Com_help				Yes	Both	No
Com_insert				Yes	Both	No
Com_insert_select				Yes	Both	No
Com_install_component				Yes	Global	No
Com_install_plugin				Yes	Both	No
Com_kill				Yes	Both	No
Com_load				Yes	Both	No
Com_lock_tables				Yes	Both	No
Com_optimize				Yes	Both	No
Com_preload_keys				Yes	Both	No
Com_prepare_sql				Yes	Both	No
Com_purge				Yes	Both	No
Com_purge_before_date				Yes	Both	No
Com_release_savepoint				Yes	Both	No
Com_rename_table				Yes	Both	No
Com_rename_user				Yes	Both	No
Com_repair				Yes	Both	No
Com_replace				Yes	Both	No
Com_replace_select				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Com_reset				Yes	Both	No
Com_resignal				Yes	Both	No
Com_revoke				Yes	Both	No
Com_revoke_all				Yes	Both	No
Com_revoke_roles				Yes	Global	No
Com_rollback				Yes	Both	No
Com_rollback_to_savepoint				Yes	Both	No
Com_savepoint				Yes	Both	No
Com_select				Yes	Both	No
Com_set_option				Yes	Both	No
Com_set_role				Yes	Global	No
Com_show_authors				Yes	Both	No
Com_show_binlog_events				Yes	Both	No
Com_show_binlogs				Yes	Both	No
Com_show_charsets				Yes	Both	No
Com_show_collations				Yes	Both	No
Com_show_contributors				Yes	Both	No
Com_show_create_db				Yes	Both	No
Com_show_create_event				Yes	Both	No
Com_show_create_func				Yes	Both	No
Com_show_create_proc				Yes	Both	No
Com_show_create_table				Yes	Both	No
Com_show_create_trigger				Yes	Both	No
Com_show_create_user				Yes	Both	No
Com_show_databases				Yes	Both	No
Com_show_engine_logs				Yes	Both	No
Com_show_engine_mutex				Yes	Both	No
Com_show_engine_status				Yes	Both	No
Com_show_errors				Yes	Both	No
Com_show_events				Yes	Both	No
Com_show_fields				Yes	Both	No
Com_show_function_code				Yes	Both	No
Com_show_function_status				Yes	Both	No
Com_show_grants				Yes	Both	No
Com_show_keys				Yes	Both	No
Com_show_master_status				Yes	Both	No
Com_show_ndb_status				Yes	Both	No
Com_show_new_master				Yes	Both	No
Com_show_open_tables				Yes	Both	No
Com_show_plugins				Yes	Both	No
Com_show_privileges				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Com_show_procedure_code				Yes	Both	No
Com_show_procedure_status				Yes	Both	No
Com_show_processlist				Yes	Both	No
Com_show_profile				Yes	Both	No
Com_show_profiles				Yes	Both	No
Com_show_relaylog_events				Yes	Both	No
Com_show_slave_hosts				Yes	Both	No
Com_show_slave_status				Yes	Both	No
Com_show_slave_status_nonblocking				Yes	Both	No
Com_show_status				Yes	Both	No
Com_show_storage_engines				Yes	Both	No
Com_show_table_status				Yes	Both	No
Com_show_tables				Yes	Both	No
Com_show_triggers				Yes	Both	No
Com_show_variables				Yes	Both	No
Com_show_warnings				Yes	Both	No
Com_shutdown				Yes	Both	No
Com_signal				Yes	Both	No
Com_slave_start				Yes	Both	No
Com_slave_stop				Yes	Both	No
Com_stmt_close				Yes	Both	No
Com_stmt_execute				Yes	Both	No
Com_stmt_fetch				Yes	Both	No
Com_stmt_prepare				Yes	Both	No
Com_stmt_reprepare				Yes	Both	No
Com_stmt_reset				Yes	Both	No
Com_stmt_send_long_data				Yes	Both	No
Com_truncate				Yes	Both	No
Com_uninstall_component				Yes	Global	No
Com_uninstall_plugin				Yes	Both	No
Com_unlock_tables				Yes	Both	No
Com_update				Yes	Both	No
Com_update_multi				Yes	Both	No
Com_xa_commit				Yes	Both	No
Com_xa_end				Yes	Both	No
Com_xa_prepare				Yes	Both	No
Com_xa_recover				Yes	Both	No
Com_xa_rollback				Yes	Both	No
Com_xa_start				Yes	Both	No
completion_type	Yes	Yes	Yes		Both	Yes
Compression				Yes	Session	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
concurrent_insert	Yes	Yes	Yes		Global	Yes
connect_timeout	Yes	Yes	Yes		Global	Yes
Connection_control_delay_generated				Yes	Global	No
connection_control_failed_connections_threshold	Yes	Yes	Yes		Global	Yes
connection_control_max_connection_delay	Yes	Yes	Yes		Global	Yes
connection_control_min_connection_delay	Yes	Yes	Yes		Global	Yes
Connection_errors_accept				Yes	Global	No
Connection_errors_internal				Yes	Global	No
Connection_errors_max_connections				Yes	Global	No
Connection_errors_peer_addr				Yes	Global	No
Connection_errors_select				Yes	Global	No
Connection_errors_tcpwrap				Yes	Global	No
Connections				Yes	Global	No
console	Yes	Yes				
core-file	Yes	Yes				
core_file			Yes		Global	No
Created_tmp_disk_tables				Yes	Both	No
Created_tmp_files				Yes	Global	No
Created_tmp_tables				Yes	Both	No
daemon_memcached_enable_binlog	Yes	Yes	Yes		Global	No
daemon_memcached_engine_lib_name	Yes	Yes	Yes		Global	No
daemon_memcached_engine_lib_path	Yes	Yes	Yes		Global	No
daemon_memcached_option	Yes	Yes	Yes		Global	No
daemon_memcached_r_batch_size	Yes	Yes	Yes		Global	No
daemon_memcached_w_batch_size	Yes	Yes	Yes		Global	No
daemonize	Yes	Yes				
datadir	Yes	Yes	Yes		Global	No
date_format			Yes		Global	No
datetime_format			Yes		Global	No
debug	Yes	Yes	Yes		Both	Yes
debug_sync			Yes		Session	Yes
debug-sync-timeout	Yes	Yes				
default_authentication_plugin	Yes	Yes	Yes		Global	No
default_password_lifetime	Yes	Yes	Yes		Global	Yes
default-storage-engine	Yes	Yes			Both	Yes
- *Variable*: default_storage_engine			Yes		Both	Yes
default-time-zone	Yes	Yes				
default_tmp_storage_engine	Yes	Yes	Yes		Both	Yes
default_week_format	Yes	Yes	Yes		Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
defaults-extra-file	Yes					
defaults-file	Yes					
defaults-group-suffix	Yes					
delay-key-write	Yes	Yes			Global	Yes
- *Variable*: delay_key_write			Yes		Global	Yes
Delayed_errors				Yes	Global	No
delayed_insert_limit	Yes	Yes	Yes		Global	Yes
Delayed_insert_threads				Yes	Global	No
delayed_insert_timeout	Yes	Yes	Yes		Global	Yes
delayed_queue_size	Yes	Yes	Yes		Global	Yes
Delayed_writes				Yes	Global	No
des-key-file	Yes	Yes				
disabled_storage_engines	Yes	Yes	Yes		Global	No
disconnect_on_expired_password	Yes	Yes	Yes		Session	No
disconnect-slave-event-count	Yes	Yes				
div_precision_increment	Yes	Yes	Yes		Both	Yes
early-plugin-load	Yes	Yes				
enable-named-pipe	Yes	Yes				
- *Variable*: named_pipe						
end_markers_in_json			Yes		Both	Yes
enforce-gtid-consistency	Yes	Yes	Yes		Global	Yes
enforce_gtid_consistency	Yes	Yes	Yes		Global	Yes
eq_range_index_dive_limit			Yes		Both	Yes
error_count			Yes		Session	No
event-scheduler	Yes	Yes			Global	Yes
- *Variable*: event_scheduler			Yes		Global	Yes
executed-gtids-compression-period	Yes	Yes				
- *Variable*: executed_gtids_compression_period						
executed_gtids_compression_period			Yes		Global	Yes
exit-info	Yes	Yes				
expire_logs_days	Yes	Yes	Yes		Global	Yes
explicit_defaults_for_timestamp	Yes	Yes	Yes		Both	Yes
external-locking	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
- *Variable*: skip_external_locking						
external_user			Yes		Session	No
federated	Yes	Yes				
flush	Yes	Yes	Yes		Global	Yes
Flush_commands				Yes	Global	No
flush_time	Yes	Yes	Yes		Global	Yes
foreign_key_checks			Yes		Both	Yes
ft_boolean_syntax	Yes	Yes	Yes		Global	Yes
ft_max_word_len	Yes	Yes	Yes		Global	No
ft_min_word_len	Yes	Yes	Yes		Global	No
ft_query_expansion_limit	Yes	Yes	Yes		Global	No
ft_stopword_file	Yes	Yes	Yes		Global	No
gdb	Yes	Yes				
general-log	Yes	Yes			Global	Yes
- *Variable*: general_log			Yes		Global	Yes
general_log_file	Yes	Yes	Yes		Global	Yes
group_concat_max_len	Yes	Yes	Yes		Both	Yes
group_replication_allow_local_disjoint_gtids_join	Yes	Yes	Yes		Global	Yes
group_replication_allow_local_lower_version_join	Yes	Yes	Yes		Global	Yes
group_replication_auto_increment_increment	Yes	Yes	Yes		Global	Yes
group_replication_bootstrap_group	Yes	Yes	Yes		Global	Yes
group_replication_components_stop_timeout	Yes	Yes	Yes		Global	Yes
group_replication_compression_threshold	Yes	Yes	Yes		Global	Yes
group_replication_enforce_update_everywhere_checks	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_applier_threshold	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_certifier_threshold	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_hold_percent	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_max_commit_quota	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_member_quota_percent	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_min_quota	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_min_recovery_quota	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_mode	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_period	Yes	Yes	Yes		Global	Yes
group_replication_flow_control_release_percent	Yes	Yes	Yes		Global	Yes
group_replication_force_members	Yes	Yes	Yes		Global	Yes
group_replication_group_name	Yes	Yes	Yes		Global	Yes
group_replication_group_seeds	Yes	Yes	Yes		Global	Yes
group_replication_gtid_assignment_block_size	Yes	Yes	Yes		Global	Yes
group_replication_ip_whitelist	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
group_replication_local_address	Yes	Yes	Yes		Global	Yes
group_replication_member_weight	Yes	Yes	Yes		Global	Yes
group_replication_poll_spin_loops	Yes	Yes	Yes		Global	Yes
group_replication_recovery_complete_at	Yes	Yes	Yes		Global	Yes
group_replication_recovery_reconnect_interval	Yes	Yes	Yes		Global	Yes
group_replication_recovery_retry_count	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_ca	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_capath	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_cert	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_cipher	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_crl	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_crlpath	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_key	Yes	Yes	Yes		Global	Yes
group_replication_recovery_ssl_verify_server_cert	Yes	Yes	Yes		Global	Yes
group_replication_recovery_use_ssl	Yes	Yes	Yes		Global	Yes
group_replication_single_primary_mode	Yes	Yes	Yes		Global	Yes
group_replication_ssl_mode	Yes	Yes	Yes		Global	Yes
group_replication_start_on_boot	Yes	Yes	Yes		Global	Yes
group_replication_transaction_size_limit	Yes	Yes	Yes		Global	Yes
group_replication_unreachable_majority_timeout	Yes	Yes	Yes		Global	Yes
gtid_executed			Yes		Varies	No
gtid-executed-compression-period	Yes	Yes				
- *Variable*: gtid_executed_compression_period						
gtid_executed_compression_period			Yes		Global	Yes
gtid-mode	Yes	Yes			Global	Yes
- *Variable*: gtid_mode			Yes		Global	Yes
gtid_mode			Yes		Global	Yes
gtid_next			Yes		Session	Yes
gtid_owned			Yes		Both	No
gtid_purged			Yes		Global	Yes
Handler_commit				Yes	Both	No
Handler_delete				Yes	Both	No
Handler_external_lock				Yes	Both	No
Handler_mrr_init				Yes	Both	No
Handler_prepare				Yes	Both	No
Handler_read_first				Yes	Both	No
Handler_read_key				Yes	Both	No
Handler_read_last				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Handler_read_next				Yes	Both	No
Handler_read_prev				Yes	Both	No
Handler_read_rnd				Yes	Both	No
Handler_read_rnd_next				Yes	Both	No
Handler_rollback				Yes	Both	No
Handler_savepoint				Yes	Both	No
Handler_savepoint_rollback				Yes	Both	No
Handler_update				Yes	Both	No
Handler_write				Yes	Both	No
have_compress			Yes		Global	No
have_crypt			Yes		Global	No
have_dynamic_loading			Yes		Global	No
have_geometry			Yes		Global	No
have_openssl			Yes		Global	No
have_profiling			Yes		Global	No
have_query_cache			Yes		Global	No
have_rtree_keys			Yes		Global	No
have_ssl			Yes		Global	No
have_statement_timeout			Yes		Global	No
have_symlink			Yes		Global	No
help	Yes	Yes				
histogram_generation_max_mem_size	Yes	Yes	Yes		Both	Yes
host_cache_size			Yes		Global	Yes
hostname			Yes		Global	No
identity			Yes		Session	Yes
ignore-builtin-innodb	Yes	Yes			Global	No
- *Variable*: ignore_builtin_innodb			Yes		Global	No
information_schema_stats	Yes	Yes	Yes		Both	Yes
init_connect	Yes	Yes	Yes		Global	Yes
init-file	Yes	Yes			Global	No
- *Variable*: init_file			Yes		Global	No
init_slave	Yes	Yes	Yes		Global	Yes
initialize	Yes	Yes				
initialize-insecure	Yes	Yes				
innodb	Yes	Yes				
innodb_adaptive_flushing	Yes	Yes	Yes		Global	Yes
innodb_adaptive_flushing_lwm	Yes	Yes	Yes		Global	Yes
innodb_adaptive_hash_index	Yes	Yes	Yes		Global	Yes
innodb_adaptive_hash_index_parts	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_adaptive_max_sleep_delay	Yes	Yes	Yes		Global	Yes
innodb_api_bk_commit_interval	Yes	Yes	Yes		Global	Yes
innodb_api_disable_rowlock	Yes	Yes	Yes		Global	No
innodb_api_enable_binlog	Yes	Yes	Yes		Global	No
innodb_api_enable_mdl	Yes	Yes	Yes		Global	No
innodb_api_trx_level	Yes	Yes	Yes		Global	Yes
innodb_autoextend_increment	Yes	Yes	Yes		Global	Yes
innodb_autoinc_lock_mode	Yes	Yes	Yes		Global	No
Innodb_available_undo_logs				Yes	Global	No
innodb_background_drop_list_empty	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_bytes_data				Yes	Global	No
Innodb_buffer_pool_bytes_dirty				Yes	Global	No
innodb_buffer_pool_chunk_size	Yes	Yes	Yes		Global	No
innodb_buffer_pool_debug	Yes	Yes	Yes		Global	No
innodb_buffer_pool_dump_at_shutdown	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_dump_now	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_dump_pct	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_dump_status				Yes	Global	No
innodb_buffer_pool_filename	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_instances	Yes	Yes	Yes		Global	No
innodb_buffer_pool_load_abort	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_load_at_startup	Yes	Yes	Yes		Global	No
innodb_buffer_pool_load_now	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_load_status				Yes	Global	No
Innodb_buffer_pool_pages_data				Yes	Global	No
Innodb_buffer_pool_pages_dirty				Yes	Global	No
Innodb_buffer_pool_pages_flushed				Yes	Global	No
Innodb_buffer_pool_pages_free				Yes	Global	No
Innodb_buffer_pool_pages_latched				Yes	Global	No
Innodb_buffer_pool_pages_misc				Yes	Global	No
Innodb_buffer_pool_pages_total				Yes	Global	No
Innodb_buffer_pool_read_ahead				Yes	Global	No
Innodb_buffer_pool_read_ahead_evicted				Yes	Global	No
Innodb_buffer_pool_read_ahead_rnd				Yes	Global	No
Innodb_buffer_pool_read_requests				Yes	Global	No
Innodb_buffer_pool_reads				Yes	Global	No
Innodb_buffer_pool_resize_status				Yes	Global	No
innodb_buffer_pool_size	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_wait_free				Yes	Global	No
Innodb_buffer_pool_write_requests				Yes	Global	No
innodb_change_buffer_max_size	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_change_buffering	Yes	Yes	Yes		Global	Yes
innodb_change_buffering_debug	Yes	Yes	Yes		Global	Yes
innodb_checkpoint_disabled	Yes	Yes	Yes		Global	Yes
innodb_checksum_algorithm	Yes	Yes	Yes		Global	Yes
innodb_cmp_per_index_enabled	Yes	Yes	Yes		Global	Yes
innodb_commit_concurrency	Yes	Yes	Yes		Global	Yes
innodb_compress_debug	Yes	Yes	Yes		Global	Yes
innodb_compression_failure_threshold_pct	Yes	Yes	Yes		Global	Yes
innodb_compression_level	Yes	Yes	Yes		Global	Yes
innodb_compression_pad_pct_max	Yes	Yes	Yes		Global	Yes
innodb_concurrency_tickets	Yes	Yes	Yes		Global	Yes
innodb_data_file_path	Yes	Yes	Yes		Global	No
Innodb_data_fsyncs				Yes	Global	No
innodb_data_home_dir	Yes	Yes	Yes		Global	No
Innodb_data_pending_fsyncs				Yes	Global	No
Innodb_data_pending_reads				Yes	Global	No
Innodb_data_pending_writes				Yes	Global	No
Innodb_data_read				Yes	Global	No
Innodb_data_reads				Yes	Global	No
Innodb_data_writes				Yes	Global	No
Innodb_data_written				Yes	Global	No
Innodb_dblwr_pages_written				Yes	Global	No
Innodb_dblwr_writes				Yes	Global	No
innodb_deadlock_detect	Yes	Yes	Yes		Global	Yes
innodb_default_row_format	Yes	Yes	Yes		Global	Yes
innodb_disable_sort_file_cache	Yes	Yes	Yes		Global	Yes
innodb_doublewrite	Yes	Yes	Yes		Global	No
innodb_fast_shutdown	Yes	Yes	Yes		Global	Yes
innodb_fil_make_page_dirty_debug	Yes	Yes	Yes		Global	Yes
innodb_file_per_table	Yes	Yes	Yes		Global	Yes
innodb_fill_factor	Yes	Yes	Yes		Global	Yes
innodb_flush_log_at_timeout			Yes		Global	Yes
innodb_flush_log_at_trx_commit	Yes	Yes	Yes		Global	Yes
innodb_flush_method	Yes	Yes	Yes		Global	No
innodb_flush_neighbors	Yes	Yes	Yes		Global	Yes
innodb_flush_sync	Yes	Yes	Yes		Global	Yes
innodb_flushing_avg_loops	Yes	Yes	Yes		Global	Yes
innodb_force_load_corrupted	Yes	Yes	Yes		Global	No
innodb_force_recovery	Yes	Yes	Yes		Global	No
innodb_ft_aux_table	Yes	Yes	Yes		Global	Yes
innodb_ft_cache_size	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_ft_enable_diag_print	Yes	Yes	Yes		Global	Yes
innodb_ft_enable_stopword	Yes	Yes	Yes		Global	Yes
innodb_ft_max_token_size	Yes	Yes	Yes		Global	No
innodb_ft_min_token_size	Yes	Yes	Yes		Global	No
innodb_ft_num_word_optimize	Yes	Yes	Yes		Global	Yes
innodb_ft_result_cache_limit	Yes	Yes	Yes		Global	Yes
innodb_ft_server_stopword_table	Yes	Yes	Yes		Global	Yes
innodb_ft_sort_pll_degree	Yes	Yes	Yes		Global	No
innodb_ft_total_cache_size	Yes	Yes	Yes		Global	No
innodb_ft_user_stopword_table	Yes	Yes	Yes		Both	Yes
Innodb_have_atomic_builtins				Yes	Global	No
innodb_io_capacity	Yes	Yes	Yes		Global	Yes
innodb_io_capacity_max	Yes	Yes	Yes		Global	Yes
innodb_limit_optimistic_insert_debug	Yes	Yes	Yes		Global	Yes
innodb_lock_wait_timeout	Yes	Yes	Yes		Both	Yes
innodb_log_buffer_size	Yes	Yes	Yes		Global	No
innodb_log_checksums	Yes	Yes	Yes		Global	Yes
innodb_log_compressed_pages	Yes	Yes	Yes		Global	Yes
innodb_log_file_size	Yes	Yes	Yes		Global	No
innodb_log_files_in_group	Yes	Yes	Yes		Global	No
innodb_log_group_home_dir	Yes	Yes	Yes		Global	No
Innodb_log_waits				Yes	Global	No
innodb_log_write_ahead_size	Yes	Yes	Yes		Global	Yes
Innodb_log_write_requests				Yes	Global	No
Innodb_log_writes				Yes	Global	No
innodb_lru_scan_depth	Yes	Yes	Yes		Global	Yes
innodb_max_dirty_pages_pct	Yes	Yes	Yes		Global	Yes
innodb_max_dirty_pages_pct_lwm	Yes	Yes	Yes		Global	Yes
innodb_max_purge_lag	Yes	Yes	Yes		Global	Yes
innodb_max_purge_lag_delay	Yes	Yes	Yes		Global	Yes
innodb_max_undo_log_size	Yes	Yes	Yes		Global	Yes
innodb_merge_threshold_set_all_debug	Yes	Yes	Yes		Global	Yes
innodb_monitor_disable	Yes	Yes	Yes		Global	Yes
innodb_monitor_enable	Yes	Yes	Yes		Global	Yes
innodb_monitor_reset	Yes	Yes	Yes		Global	Yes
innodb_monitor_reset_all	Yes	Yes	Yes		Global	Yes
Innodb_num_open_files				Yes	Global	No
innodb_numa_interleave	Yes	Yes	Yes		Global	No
innodb_old_blocks_pct	Yes	Yes	Yes		Global	Yes
innodb_old_blocks_time	Yes	Yes	Yes		Global	Yes
innodb_online_alter_log_max_size	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_open_files	Yes	Yes	Yes		Global	No
innodb_optimize_fulltext_only	Yes	Yes	Yes		Global	Yes
Innodb_os_log_fsyncs				Yes	Global	No
Innodb_os_log_pending_fsyncs				Yes	Global	No
Innodb_os_log_pending_writes				Yes	Global	No
Innodb_os_log_written				Yes	Global	No
innodb_page_cleaners	Yes	Yes	Yes		Global	No
Innodb_page_size				Yes	Global	No
innodb_page_size	Yes	Yes	Yes		Global	No
Innodb_pages_created				Yes	Global	No
Innodb_pages_read				Yes	Global	No
Innodb_pages_written				Yes	Global	No
innodb_print_all_deadlocks	Yes	Yes	Yes		Global	Yes
innodb_purge_batch_size	Yes	Yes	Yes		Global	Yes
innodb_purge_rseg_truncate_frequency	Yes	Yes	Yes		Global	Yes
innodb_purge_threads	Yes	Yes	Yes		Global	No
innodb_random_read_ahead	Yes	Yes	Yes		Global	Yes
innodb_read_ahead_threshold	Yes	Yes	Yes		Global	Yes
innodb_read_io_threads	Yes	Yes	Yes		Global	No
innodb_read_only	Yes	Yes	Yes		Global	No
innodb_redo_log_encrypt	Yes	Yes	Yes		Global	Yes
innodb_replication_delay	Yes	Yes	Yes		Global	Yes
innodb_rollback_on_timeout	Yes	Yes	Yes		Global	No
innodb_rollback_segments	Yes	Yes	Yes		Global	Yes
Innodb_row_lock_current_waits				Yes	Global	No
Innodb_row_lock_time				Yes	Global	No
Innodb_row_lock_time_avg				Yes	Global	No
Innodb_row_lock_time_max				Yes	Global	No
Innodb_row_lock_waits				Yes	Global	No
Innodb_rows_deleted				Yes	Global	No
Innodb_rows_inserted				Yes	Global	No
Innodb_rows_read				Yes	Global	No
Innodb_rows_updated				Yes	Global	No
innodb_saved_page_number_debug	Yes	Yes	Yes		Global	Yes
innodb_scan_directories	Yes	Yes	Yes		Global	No
innodb_sort_buffer_size	Yes	Yes	Yes		Global	No
innodb_spin_wait_delay	Yes	Yes	Yes		Global	Yes
innodb_stats_auto_recalc	Yes	Yes	Yes		Global	Yes
innodb_stats_include_delete_marked	Yes	Yes	Yes		Global	Yes
innodb_stats_method	Yes	Yes	Yes		Global	Yes
innodb_stats_on_metadata	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_stats_persistent	Yes	Yes	Yes		Global	Yes
innodb_stats_persistent_sample_pages	Yes	Yes	Yes		Global	Yes
innodb_stats_transient_sample_pages	Yes	Yes	Yes		Global	Yes
innodb-status-file	Yes	Yes				
innodb_status_output	Yes	Yes	Yes		Global	Yes
innodb_status_output_locks	Yes	Yes	Yes		Global	Yes
innodb_strict_mode	Yes	Yes	Yes		Both	Yes
innodb_sync_array_size	Yes	Yes	Yes		Global	No
innodb_sync_debug	Yes	Yes	Yes		Global	No
innodb_sync_spin_loops	Yes	Yes	Yes		Global	Yes
innodb_table_locks	Yes	Yes	Yes		Both	Yes
innodb_temp_data_file_path	Yes	Yes	Yes		Global	No
innodb_thread_concurrency	Yes	Yes	Yes		Global	Yes
innodb_thread_sleep_delay	Yes	Yes	Yes		Global	Yes
innodb_tmpdir	Yes	Yes	Yes		Both	Yes
Innodb_truncated_status_writes				Yes	Global	No
innodb_trx_purge_view_update_only_debug	Yes	Yes	Yes		Global	Yes
innodb_trx_rseg_n_slots_debug	Yes	Yes	Yes		Global	Yes
innodb_undo_directory	Yes	Yes	Yes		Global	No
innodb_undo_log_encrypt	Yes	Yes	Yes		Global	Yes
innodb_undo_log_truncate	Yes	Yes	Yes		Global	Yes
innodb_undo_logs	Yes	Yes	Yes		Global	Yes
innodb_undo_tablespaces	Yes	Yes	Yes		Global	Varies
innodb_use_native_aio	Yes	Yes	Yes		Global	No
innodb_version			Yes		Global	No
innodb_write_io_threads	Yes	Yes	Yes		Global	No
insert_id			Yes		Session	Yes
install	Yes					
install-manual	Yes					
interactive_timeout	Yes	Yes	Yes		Both	Yes
internal_tmp_disk_storage_engine	Yes	Yes	Yes		Global	Yes
internal_tmp_mem_storage_engine	Yes	Yes	Yes		Session	Yes
join_buffer_size	Yes	Yes	Yes		Both	Yes
keep_files_on_create	Yes	Yes	Yes		Both	Yes
Key_blocks_not_flushed				Yes	Global	No
Key_blocks_unused				Yes	Global	No
Key_blocks_used				Yes	Global	No
key_buffer_size	Yes	Yes	Yes		Global	Yes
key_cache_age_threshold	Yes	Yes	Yes		Global	Yes
key_cache_block_size	Yes	Yes	Yes		Global	Yes
key_cache_division_limit	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Key_read_requests				Yes	Global	No
Key_reads				Yes	Global	No
Key_write_requests				Yes	Global	No
Key_writes				Yes	Global	No
keyring_file_data	Yes	Yes	Yes		Global	Yes
language	Yes	Yes	Yes		Global	No
large_files_support			Yes		Global	No
large_page_size			Yes		Global	No
large-pages	Yes	Yes			Global	No
- *Variable*: large_pages			Yes		Global	No
last_insert_id			Yes		Session	Yes
Last_query_cost				Yes	Session	No
Last_query_partial_plans				Yes	Session	No
lc-messages	Yes	Yes			Both	Yes
- *Variable*: lc_messages			Yes		Both	Yes
lc-messages-dir	Yes	Yes			Global	No
- *Variable*: lc_messages_dir			Yes		Global	No
lc_time_names			Yes		Both	Yes
license			Yes		Global	No
local_infile			Yes		Global	Yes
local-service	Yes					
lock_wait_timeout	Yes	Yes	Yes		Both	Yes
Locked_connects				Yes	Global	No
locked_in_memory			Yes		Global	No
log-bin	Yes	Yes	Yes		Global	No
log_bin			Yes		Global	No
log_bin_basename			Yes		Global	No
log-bin-index	Yes	Yes				
log_bin_index			Yes		Global	No
log-bin-trust-function-creators	Yes	Yes			Global	Yes
- *Variable*: log_bin_trust_function_creators			Yes		Global	Yes
log-bin-use-v1-row-events	Yes	Yes			Global	No
- *Variable*: log_bin_use_v1_row_events			Yes		Global	No
log_bin_use_v1_row_events	Yes	Yes	Yes		Global	No
log_builtin_as_identified_by_password	Yes	Yes	Yes		Global	Yes
log-error	Yes	Yes			Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
- *Variable*: log_error			Yes		Global	No
log_error_verbosity	Yes	Yes	Yes		Global	Yes
log-isam	Yes	Yes				
log-output	Yes	Yes			Global	Yes
- *Variable*: log_output			Yes		Global	Yes
log-queries-not-using-indexes	Yes	Yes			Global	Yes
- *Variable*: log_queries_not_using_indexes			Yes		Global	Yes
log-raw	Yes	Yes				
log-short-format	Yes	Yes				
log-slave-updates	Yes	Yes			Global	No
- *Variable*: log_slave_updates			Yes		Global	No
log_slave_updates	Yes	Yes	Yes		Global	No
log_slow_admin_statements			Yes		Global	Yes
log_slow_slave_statements			Yes		Global	Yes
log_statements_unsafe_for_binlog			Yes		Global	Yes
log_syslog	Yes	Yes	Yes		Global	Yes
log_syslog_facility	Yes	Yes	Yes		Global	Yes
log_syslog_include_pid	Yes	Yes	Yes		Global	Yes
log_syslog_tag	Yes	Yes	Yes		Global	Yes
log-tc	Yes	Yes				
log-tc-size	Yes	Yes				
log_throttle_queries_not_using_indexes			Yes		Global	Yes
log_timestamps	Yes	Yes	Yes		Global	Yes
log-warnings	Yes	Yes			Global	Yes
- *Variable*: log_warnings			Yes		Global	Yes
long_query_time	Yes	Yes	Yes		Both	Yes
low-priority-updates	Yes	Yes			Both	Yes
- *Variable*: low_priority_updates			Yes		Both	Yes
lower_case_file_system			Yes		Global	No
lower_case_table_names	Yes	Yes	Yes		Global	No
master-info-file	Yes	Yes				
master-info-repository	Yes	Yes				
- *Variable*: master_info_repository						

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
master_info_repository	Yes	Yes	Yes		Global	Yes
master-retry-count	Yes	Yes				
master-verify-checksum	Yes	Yes				
- *Variable*: master_verify_checksum						
master_verify_checksum			Yes		Global	Yes
max_allowed_packet	Yes	Yes	Yes		Both	Yes
max_binlog_cache_size	Yes	Yes	Yes		Global	Yes
max-binlog-dump-events	Yes	Yes				
max_binlog_size	Yes	Yes	Yes		Global	Yes
max_binlog_stmt_cache_size	Yes	Yes	Yes		Global	Yes
max_connect_errors	Yes	Yes	Yes		Global	Yes
max_connections	Yes	Yes	Yes		Global	Yes
max_delayed_threads	Yes	Yes	Yes		Both	Yes
max_digest_length	Yes	Yes	Yes		Global	No
max_error_count	Yes	Yes	Yes		Both	Yes
max_execution_time	Yes	Yes	Yes		Both	Yes
Max_execution_time_exceeded				Yes	Both	No
Max_execution_time_set				Yes	Both	No
Max_execution_time_set_failed				Yes	Both	No
max_heap_table_size	Yes	Yes	Yes		Both	Yes
max_insert_delayed_threads			Yes		Both	Yes
max_join_size	Yes	Yes	Yes		Both	Yes
max_length_for_sort_data	Yes	Yes	Yes		Both	Yes
max_points_in_geometry	Yes	Yes	Yes		Global	Yes
max_prepared_stmt_count	Yes	Yes	Yes		Global	Yes
max_relay_log_size	Yes	Yes	Yes		Global	Yes
max_seeks_for_key	Yes	Yes	Yes		Both	Yes
max_sort_length	Yes	Yes	Yes		Both	Yes
max_sp_recursion_depth	Yes	Yes	Yes		Both	Yes
max_tmp_tables			Yes		Both	Yes
Max_used_connections				Yes	Global	No
Max_used_connections_time				Yes	Global	No
max_user_connections	Yes	Yes	Yes		Both	Yes
max_write_lock_count	Yes	Yes	Yes		Global	Yes
mecab_charset				Yes	Global	No
mecab_rc_file	Yes	Yes	Yes		Global	No
memlock	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
- *Variable*: locked_in_memory						
metadata_locks_cache_size			Yes		Global	No
metadata_locks_hash_instances			Yes		Global	No
min-examined-row-limit	Yes	Yes	Yes		Both	Yes
myisam-block-size	Yes	Yes				
myisam_data_pointer_size	Yes	Yes	Yes		Global	Yes
myisam_max_sort_file_size	Yes	Yes	Yes		Global	Yes
myisam_mmap_size	Yes	Yes	Yes		Global	No
myisam-recover-options	Yes	Yes				
- *Variable*: myisam_recover_options						
myisam_recover_options			Yes		Global	No
myisam_repair_threads	Yes	Yes	Yes		Both	Yes
myisam_sort_buffer_size	Yes	Yes	Yes		Both	Yes
myisam_stats_method	Yes	Yes	Yes		Both	Yes
myisam_use_mmap	Yes	Yes	Yes		Global	Yes
mysql_native_password_proxy_users	Yes		Yes		Global	Yes
Mysqlx_address				Yes	Global	No
mysqlx_bind_address	Yes	Yes	Yes		Global	No
Mysqlx_bytes_received				Yes	Both	No
Mysqlx_bytes_sent				Yes	Both	No
mysqlx_connect_timeout	Yes	Yes	Yes		Global	Yes
Mysqlx_connection_accept_errors				Yes	Both	No
Mysqlx_connection_errors				Yes	Both	No
Mysqlx_connections_accepted				Yes	Global	No
Mysqlx_connections_closed				Yes	Global	No
Mysqlx_connections_rejected				Yes	Global	No
Mysqlx_crud_create_view				Yes	Both	No
Mysqlx_crud_delete				Yes	Both	No
Mysqlx_crud_drop_view				Yes	Both	No
Mysqlx_crud_find				Yes	Both	No
Mysqlx_crud_insert				Yes	Both	No
Mysqlx_crud_modify_view				Yes	Both	No
Mysqlx_crud_update				Yes	Both	No
Mysqlx_errors_sent				Yes	Both	No
Mysqlx_errors_unknown_message_type				Yes	Both	No
Mysqlx_expect_close				Yes	Both	No
Mysqlx_expect_open				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
mysqlx_idle_worker_thread_timeout	Yes	Yes				
Mysqlx_init_error				Yes	Both	No
mysqlx_max_allowed_packet	Yes	Yes				
mysqlx_max_connections	Yes	Yes	Yes		Global	Yes
mysqlx_min_worker_threads	Yes	Yes				
Mysqlx_notice_other_sent				Yes	Both	No
Mysqlx_notice_warning_sent				Yes	Both	No
Mysqlx_port				Yes	Global	No
mysqlx_port	Yes	Yes	Yes		Global	No
mysqlx_port_open_timeout	Yes	Yes	Yes		Global	No
Mysqlx_rows_sent				Yes	Both	No
Mysqlx_sessions				Yes	Global	No
Mysqlx_sessions_accepted				Yes	Global	No
Mysqlx_sessions_closed				Yes	Global	No
Mysqlx_sessions_fatal_error				Yes	Global	No
Mysqlx_sessions_killed				Yes	Global	No
Mysqlx_sessions_rejected				Yes	Global	No
Mysqlx_socket				Yes	Global	No
mysqlx-socket	Yes	Yes	Yes		Global	No
mysqlx_socket	Yes	Yes	Yes		Global	No
mysqlx_ssl	Yes	Yes	Yes		Global	No
Mysqlx_ssl_accept_renegotiates				Yes	Global	No
Mysqlx_ssl_accepts				Yes	Global	No
Mysqlx_ssl_active				Yes	Both	No
mysqlx_ssl_ca	Yes	Yes	Yes		Global	No
mysqlx_ssl_capath	Yes	Yes	Yes		Global	No
mysqlx_ssl_cert	Yes	Yes	Yes		Global	No
Mysqlx_ssl_cipher				Yes	Both	No
mysqlx_ssl_cipher	Yes	Yes				
Mysqlx_ssl_cipher_list				Yes	Both	No
mysqlx_ssl_crl	Yes	Yes	Yes		Global	No
mysqlx_ssl_crlpath	Yes	Yes	Yes		Global	No
Mysqlx_ssl_ctx_verify_depth				Yes	Both	No
Mysqlx_ssl_ctx_verify_mode				Yes	Both	No
Mysqlx_ssl_finished_accepts				Yes	Global	No
mysqlx_ssl_key	Yes	Yes	Yes		Global	No
Mysqlx_ssl_server_not_after				Yes	Global	No
Mysqlx_ssl_server_not_before				Yes	Global	No
Mysqlx_ssl_verify_depth				Yes	Global	No
Mysqlx_ssl_verify_mode				Yes	Global	No
Mysqlx_ssl_version				Yes	Both	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Mysqlx_stmt_create_collection				Yes	Both	No
Mysqlx_stmt_create_collection_index				Yes	Both	No
Mysqlx_stmt_disable_notices				Yes	Both	No
Mysqlx_stmt_drop_collection				Yes	Both	No
Mysqlx_stmt_drop_collection_index				Yes	Both	No
Mysqlx_stmt_enable_notices				Yes	Both	No
Mysqlx_stmt_ensure_collection				Yes	Both	No
Mysqlx_stmt_execute_mysqlx				Yes	Both	No
Mysqlx_stmt_execute_sql				Yes	Both	No
Mysqlx_stmt_execute_xplugin				Yes	Both	No
Mysqlx_stmt_kill_client				Yes	Both	No
Mysqlx_stmt_list_clients				Yes	Both	No
Mysqlx_stmt_list_notices				Yes	Both	No
Mysqlx_stmt_list_objects				Yes	Both	No
Mysqlx_stmt_ping				Yes	Both	No
Mysqlx_worker_threads				Yes	Global	No
Mysqlx_worker_threads_active				Yes	Global	No
named_pipe			Yes		Global	No
Ndb_api_bytes_received_count				Yes	Global	No
Ndb_api_bytes_received_count_session				Yes	Session	No
Ndb_api_bytes_received_count_slave				Yes	Global	No
Ndb_api_bytes_sent_count				Yes	Global	No
Ndb_api_bytes_sent_count_slave				Yes	Global	No
Ndb_api_event_bytes_count_injector				Yes	Global	No
Ndb_api_event_data_count_injector				Yes	Global	No
Ndb_api_event_nondata_count_injector				Yes	Global	No
Ndb_api_pk_op_count				Yes	Global	No
Ndb_api_pk_op_count_session				Yes	Session	No
Ndb_api_pk_op_count_slave				Yes	Global	No
Ndb_api_pruned_scan_count				Yes	Global	No
Ndb_api_pruned_scan_count_session				Yes	Session	No
Ndb_api_range_scan_count_slave				Yes	Global	No
Ndb_api_read_row_count				Yes	Global	No
Ndb_api_read_row_count_session				Yes	Session	No
Ndb_api_scan_batch_count_slave				Yes	Global	No
Ndb_api_table_scan_count				Yes	Global	No
Ndb_api_table_scan_count_session				Yes	Session	No
Ndb_api_trans_abort_count				Yes	Global	No
Ndb_api_trans_abort_count_session				Yes	Session	No
Ndb_api_trans_abort_count_slave				Yes	Global	No
Ndb_api_trans_close_count				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Ndb_api_trans_close_count_session				Yes	Session	No
Ndb_api_trans_close_count_slave				Yes	Global	No
Ndb_api_trans_commit_count				Yes	Global	No
Ndb_api_trans_commit_count_session				Yes	Session	No
Ndb_api_trans_commit_count_slave				Yes	Global	No
Ndb_api_trans_local_read_row_count_slave				Yes	Global	No
Ndb_api_trans_start_count				Yes	Global	No
Ndb_api_trans_start_count_session				Yes	Session	No
Ndb_api_trans_start_count_slave				Yes	Global	No
Ndb_api_uk_op_count				Yes	Global	No
Ndb_api_uk_op_count_slave				Yes	Global	No
Ndb_api_wait_exec_complete_count				Yes	Global	No
Ndb_api_wait_exec_complete_count_session				Yes	Session	No
Ndb_api_wait_exec_complete_count_slave				Yes	Global	No
Ndb_api_wait_meta_request_count				Yes	Global	No
Ndb_api_wait_meta_request_count_session				Yes	Session	No
Ndb_api_wait_nanos_count				Yes	Global	No
Ndb_api_wait_nanos_count_session				Yes	Session	No
Ndb_api_wait_nanos_count_slave				Yes	Global	No
Ndb_api_wait_scan_result_count				Yes	Global	No
Ndb_api_wait_scan_result_count_session				Yes	Session	No
Ndb_api_wait_scan_result_count_slave				Yes	Global	No
ndb-batch-size	Yes	Yes	Yes		Global	No
ndb-blob-write-batch-bytes	Yes	Yes	Yes		Both	Yes
ndb-cluster-connection-pool	Yes	Yes	Yes		Global	No
ndb-cluster-connection-pool-nodeids	Yes	Yes	Yes		Global	No
Ndb_cluster_node_id				Yes	Both	No
Ndb_config_from_host				Yes	Both	No
Ndb_config_from_port				Yes	Both	No
Ndb_conflict_fn_epoch_trans				Yes	Global	No
Ndb_conflict_fn_max				Yes	Global	No
Ndb_conflict_fn_old				Yes	Global	No
Ndb_conflict_trans_detect_iter_count				Yes	Global	No
Ndb_conflict_trans_row_reject_count				Yes	Global	No
ndb-connectstring	Yes	Yes				
ndb-deferred-constraints	Yes	Yes			Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
- *Variable*: ndb_deferred_constraints			Yes		Both	Yes
ndb_deferred_constraints	Yes	Yes	Yes		Both	Yes
ndb-distribution	Yes	Yes			Global	Yes
- *Variable*: ndb_distribution			Yes		Global	Yes
ndb_distribution	Yes	Yes	Yes		Global	Yes
ndb_eventbuffer_free_percent	Yes	Yes	Yes		Global	Yes
ndb_eventbuffer_max_alloc	Yes	Yes	Yes		Global	Yes
ndb_force_send	Yes	Yes	Yes		Both	Yes
ndb_index_stat_enable	Yes	Yes	Yes		Both	Yes
ndb_index_stat_option	Yes	Yes	Yes		Both	Yes
ndb_join_pushdown			Yes		Both	Yes
Ndb_last_commit_epoch_server				Yes	Global	No
Ndb_last_commit_epoch_session				Yes	Session	No
ndb-log-apply-status	Yes	Yes			Global	No
- *Variable*: ndb_log_apply_status			Yes		Global	No
ndb_log_apply_status	Yes	Yes	Yes		Global	No
ndb_log_binlog_index	Yes		Yes		Global	Yes
ndb-log-empty-epochs	Yes	Yes	Yes		Global	Yes
ndb-log-empty-update	Yes	Yes	Yes		Global	Yes
ndb-log-transaction-id	Yes	Yes			Global	No
- *Variable*: ndb_log_transaction_id			Yes		Global	No
ndb_log_updated_only	Yes	Yes	Yes		Global	Yes
ndb-mgmd-host	Yes	Yes				
Ndb_number_of_data_nodes				Yes	Global	No
ndb_optimization_delay			Yes		Global	Yes
ndb_optimized_node_selection	Yes	Yes	Yes		Global	No
Ndb_pushed_queries_defined				Yes	Global	No
Ndb_pushed_queries_executed				Yes	Global	No
ndb_recv_thread_activation_threshold						
ndb_recv_thread_cpu_mask			Yes		Global	Yes
ndb_report_thresh_binlog_epoch_slip	Yes	Yes	Yes		Global	Yes
ndb_report_thresh_binlog_mem_usage	Yes	Yes	Yes		Global	Yes
Ndb_scan_count				Yes	Global	No
ndb_show_foreign_key_mock_tables	Yes	Yes	Yes		Global	Yes
Ndb_slave_max_replicated_epoch			Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
ndb_table_no_logging			Yes		Session	Yes
ndb-transid-mysql-connection-map	Yes					
ndb_use_transactions	Yes	Yes	Yes		Both	Yes
ndb_version			Yes		Global	No
ndb_version_string			Yes		Global	No
ndb-wait-setup	Yes	Yes	Yes		Global	No
ndbinfo_database			Yes		Global	No
ndbinfo_max_rows	Yes		Yes		Both	Yes
ndbinfo_show_hidden	Yes		Yes		Both	Yes
ndbinfo_version			Yes		Global	No
net_buffer_length	Yes	Yes	Yes		Both	Yes
net_read_timeout	Yes	Yes	Yes		Both	Yes
net_retry_count	Yes	Yes	Yes		Both	Yes
net_write_timeout	Yes	Yes	Yes		Both	Yes
new	Yes	Yes	Yes		Both	Yes
ngram_token_size	Yes	Yes	Yes		Global	No
no-defaults	Yes					
Not_flushed_delayed_rows				Yes	Global	No
offline_mode	Yes	Yes	Yes		Global	Yes
old	Yes	Yes	Yes		Global	No
old-alter-table	Yes	Yes			Both	Yes
- *Variable*: old_alter_table			Yes		Both	Yes
old_passwords			Yes		Both	Yes
old-style-user-limits	Yes	Yes				
Ongoing_anonymous_gtid_violating_transaction_count				Yes	Global	No
Ongoing_anonymous_transaction_count				Yes	Global	No
Ongoing_automatic_gtid_violating_transaction_count				Yes	Global	No
Open_files				Yes	Global	No
open-files-limit	Yes	Yes			Global	No
- *Variable*: open_files_limit			Yes		Global	No
Open_streams				Yes	Global	No
Open_table_definitions				Yes	Global	No
Open_tables				Yes	Both	No
Opened_files				Yes	Global	No
Opened_table_definitions				Yes	Both	No
Opened_tables				Yes	Both	No
optimizer_prune_level	Yes	Yes	Yes		Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
optimizer_search_depth	Yes	Yes	Yes		Both	Yes
optimizer_switch	Yes	Yes	Yes		Both	Yes
optimizer_trace			Yes		Both	Yes
optimizer_trace_features			Yes		Both	Yes
optimizer_trace_limit			Yes		Both	Yes
optimizer_trace_max_mem_size			Yes		Both	Yes
optimizer_trace_offset			Yes		Both	Yes
original_commit_timestamp				Yes	Session	Yes
parser_max_mem_size	Yes	Yes	Yes		Both	Yes
performance_schema	Yes	Yes	Yes		Global	No
Performance_schema_accounts_lost				Yes	Global	No
performance_schema_accounts_size	Yes	Yes	Yes		Global	No
Performance_schema_cond_classes_lost				Yes	Global	No
Performance_schema_cond_instances_lost				Yes	Global	No
performance-schema-consumer-events-stages-current	Yes	Yes				
performance-schema-consumer-events-stages-history	Yes	Yes				
performance-schema-consumer-events-stages-history-long	Yes	Yes				
performance-schema-consumer-events-statements-current	Yes	Yes				
performance-schema-consumer-events-statements-history	Yes	Yes				
performance-schema-consumer-events-statements-history-long	Yes	Yes				
performance-schema-	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
consumer-events-transactions-current						
performance-schema-consumer-events-transactions-history	Yes	Yes				
performance-schema-consumer-events-transactions-history-long	Yes	Yes				
performance-schema-consumer-events-waits-current	Yes	Yes				
performance-schema-consumer-events-waits-history	Yes	Yes				
performance-schema-consumer-events-waits-history-long	Yes	Yes				
performance-schema-consumer-global-instrumentation	Yes	Yes				
performance-schema-consumer-statements-digest	Yes	Yes				
performance-schema-consumer-thread-instrumentation	Yes	Yes				
Performance_schema_digest_lost				Yes	Global	No
performance_schema_digests_size	Yes	Yes	Yes		Global	No
performance_schema_error_size	Yes	Yes	Yes		Global	No
performance_schema_events_stages_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_stages_history_size	Yes	Yes	Yes		Global	No
performance_schema_events_statements_history_long_size	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
performance_schema_events_statements_history_size	Yes	Yes	Yes		Global	No
performance_schema_events_transactions_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_transactions_history_size	Yes	Yes	Yes		Global	No
performance_schema_events_waits_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_waits_history_size	Yes	Yes	Yes		Global	No
Performance_schema_file_classes_lost				Yes	Global	No
Performance_schema_file_handles_lost				Yes	Global	No
Performance_schema_file_instances_lost				Yes	Global	No
Performance_schema_hosts_lost				Yes	Global	No
performance_schema_hosts_size	Yes	Yes	Yes		Global	No
Performance_schema_index_stat_lost				Yes	Global	No
performance-schema-instrument	Yes	Yes				
Performance_schema_locker_lost				Yes	Global	No
performance_schema_max_cond_classes	Yes	Yes	Yes		Global	No
performance_schema_max_cond_instances	Yes	Yes	Yes		Global	No
performance_schema_max_digest_length	Yes	Yes	Yes		Global	No
performance_schema_max_file_classes	Yes	Yes	Yes		Global	No
performance_schema_max_file_handles	Yes	Yes	Yes		Global	No
performance_schema_max_file_instances	Yes	Yes	Yes		Global	No
performance_schema_max_index_stat	Yes	Yes	Yes		Global	No
performance_schema_max_memory_classes	Yes	Yes	Yes		Global	No
performance_schema_max_metadata_locks	Yes	Yes	Yes		Global	No
performance_schema_max_mutex_classes	Yes	Yes	Yes		Global	No
performance_schema_max_mutex_instances	Yes	Yes	Yes		Global	No
performance_schema_max_prepared_statements_instances	Yes	Yes	Yes		Global	No
performance_schema_max_program_instances	Yes	Yes	Yes		Global	No
performance_schema_max_rwlock_classes	Yes	Yes	Yes		Global	No
performance_schema_max_rwlock_instances	Yes	Yes	Yes		Global	No
performance_schema_max_socket_classes	Yes	Yes	Yes		Global	No
performance_schema_max_socket_instances	Yes	Yes	Yes		Global	No
performance_schema_max_sql_text_length	Yes	Yes	Yes		Global	No
performance_schema_max_stage_classes	Yes	Yes	Yes		Global	No
performance_schema_max_statement_classes	Yes	Yes	Yes		Global	No
performance_schema_max_statement_stack	Yes	Yes	Yes		Global	No
performance_schema_max_table_handles	Yes	Yes	Yes		Global	No
performance_schema_max_table_instances	Yes	Yes	Yes		Global	No
performance_schema_max_table_lock_stat	Yes	Yes	Yes		Global	No
performance_schema_max_thread_classes	Yes	Yes	Yes		Global	No
performance_schema_max_thread_instances	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Performance_schema_memory_classes_lost				Yes	Global	No
Performance_schema_metadata_lock_lost				Yes	Global	No
Performance_schema_mutex_classes_lost				Yes	Global	No
Performance_schema_mutex_instances_lost				Yes	Global	No
Performance_schema_nested_statement_lost				Yes	Global	No
Performance_schema_prepared_statements_lost				Yes	Global	No
Performance_schema_program_lost				Yes	Global	No
Performance_schema_rwlock_classes_lost				Yes	Global	No
Performance_schema_rwlock_instances_lost				Yes	Global	No
Performance_schema_session_connect_attrs_longest_seen				Yes	Global	No
Performance_schema_session_connect_attrs_lost				Yes	Global	No
performance_schema_session_connect_attrs_size	Yes	Yes	Yes		Global	No
performance_schema_setup_actors_size	Yes	Yes	Yes		Global	No
performance_schema_setup_objects_size	Yes	Yes	Yes		Global	No
Performance_schema_socket_classes_lost				Yes	Global	No
Performance_schema_socket_instances_lost				Yes	Global	No
Performance_schema_stage_classes_lost				Yes	Global	No
Performance_schema_statement_classes_lost				Yes	Global	No
Performance_schema_table_handles_lost				Yes	Global	No
Performance_schema_table_instances_lost				Yes	Global	No
Performance_schema_table_lock_stat_lost				Yes	Global	No
Performance_schema_thread_classes_lost				Yes	Global	No
Performance_schema_thread_instances_lost				Yes	Global	No
Performance_schema_users_lost				Yes	Global	No
performance_schema_users_size	Yes	Yes	Yes		Global	No
persisted_globals_load	Yes	Yes	Yes		Global	No
pid-file	Yes	Yes			Global	No
- *Variable*: pid_file			Yes		Global	No
plugin	Yes	Yes				
plugin_dir	Yes	Yes	Yes		Global	No
plugin-load	Yes	Yes				
plugin-load-add	Yes	Yes				
port	Yes	Yes	Yes		Global	No
port-open-timeout	Yes	Yes				
preload_buffer_size	Yes	Yes	Yes		Both	Yes
Prepared_stmt_count				Yes	Global	No
print-defaults	Yes					
profiling			Yes		Both	Yes
profiling_history_size	Yes	Yes	Yes		Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
protocol_version			Yes		Global	No
proxy_user			Yes		Session	No
pseudo_slave_mode			Yes		Session	Yes
pseudo_thread_id			Yes		Session	Yes
Qcache_free_blocks				Yes	Global	No
Qcache_free_memory				Yes	Global	No
Qcache_hits				Yes	Global	No
Qcache_inserts				Yes	Global	No
Qcache_lowmem_prunes				Yes	Global	No
Qcache_not_cached				Yes	Global	No
Qcache_queries_in_cache				Yes	Global	No
Qcache_total_blocks				Yes	Global	No
Queries				Yes	Both	No
query_alloc_block_size	Yes	Yes	Yes		Both	Yes
query_cache_limit	Yes	Yes	Yes		Global	Yes
query_cache_min_res_unit	Yes	Yes	Yes		Global	Yes
query_cache_size	Yes	Yes	Yes		Global	Yes
query_cache_type	Yes	Yes	Yes		Both	Yes
query_cache_wlock_invalidate	Yes	Yes	Yes		Both	Yes
query_prealloc_size	Yes	Yes	Yes		Both	Yes
Questions				Yes	Both	No
rand_seed1			Yes		Session	Yes
rand_seed2			Yes		Session	Yes
range_alloc_block_size	Yes	Yes	Yes		Both	Yes
range_optimizer_max_mem_size	Yes	Yes	Yes		Both	Yes
rbr_exec_mode			Yes		Session	Yes
read_buffer_size	Yes	Yes	Yes		Both	Yes
read_only	Yes	Yes	Yes		Global	Yes
read_rnd_buffer_size	Yes	Yes	Yes		Both	Yes
relay-log	Yes	Yes			Global	No
- *Variable*: relay_log			Yes		Global	No
relay_log_basename			Yes		Global	No
relay-log-index	Yes	Yes			Global	No
- *Variable*: relay_log_index			Yes		Global	No
relay_log_index	Yes	Yes	Yes		Global	No
relay-log-info-file	Yes	Yes				
- *Variable*: relay_log_info_file						
relay_log_info_file	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
relay-log-info-repository	Yes	Yes				
- *Variable*: relay_log_info_repository						
relay_log_info_repository			Yes		Global	Yes
relay_log_purge	Yes	Yes	Yes		Global	Yes
relay-log-recovery	Yes	Yes				
- *Variable*: relay_log_recovery						
relay_log_recovery	Yes	Yes	Yes		Global	No
relay_log_space_limit	Yes	Yes	Yes		Global	No
remove	Yes					
replicate-do-db	Yes	Yes				
replicate-do-table	Yes	Yes				
replicate-ignore-db	Yes	Yes				
replicate-ignore-table	Yes	Yes				
replicate-rewrite-db	Yes	Yes				
replicate-same-server-id	Yes	Yes				
replicate-wild-do-table	Yes	Yes				
replicate-wild-ignore-table	Yes	Yes				
report-host	Yes	Yes			Global	No
- *Variable*: report_host			Yes		Global	No
report-password	Yes	Yes			Global	No
- *Variable*: report_password			Yes		Global	No
report-port	Yes	Yes			Global	No
- *Variable*: report_port			Yes		Global	No
report-user	Yes	Yes			Global	No
- *Variable*: report_user			Yes		Global	No
require_secure_transport	Yes	Yes	Yes		Global	Yes
rewriter_enabled			Yes		Global	Yes
Rewriter_number_loaded_rules				Yes	Global	No
Rewriter_number_reloads				Yes	Global	No
Rewriter_number_rewritten_queries				Yes	Global	No
Rewriter_reload_error				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
rewriter_verbose			Yes		Global	Yes
Rpl_semi_sync_master_clients				Yes	Global	No
rpl_semi_sync_master_enabled			Yes		Global	Yes
Rpl_semi_sync_master_net_avg_wait_time				Yes	Global	No
Rpl_semi_sync_master_net_wait_time				Yes	Global	No
Rpl_semi_sync_master_net_waits				Yes	Global	No
Rpl_semi_sync_master_no_times				Yes	Global	No
Rpl_semi_sync_master_no_tx				Yes	Global	No
Rpl_semi_sync_master_status				Yes	Global	No
Rpl_semi_sync_master_timefunc_failures				Yes	Global	No
rpl_semi_sync_master_timeout			Yes		Global	Yes
rpl_semi_sync_master_trace_level			Yes		Global	Yes
Rpl_semi_sync_master_tx_avg_wait_time				Yes	Global	No
Rpl_semi_sync_master_tx_wait_time				Yes	Global	No
Rpl_semi_sync_master_tx_waits				Yes	Global	No
rpl_semi_sync_master_wait_for_slave_count			Yes		Global	Yes
rpl_semi_sync_master_wait_no_slave			Yes		Global	Yes
rpl_semi_sync_master_wait_point			Yes		Global	Yes
Rpl_semi_sync_master_wait_pos_backtraverse				Yes	Global	No
Rpl_semi_sync_master_wait_sessions				Yes	Global	No
Rpl_semi_sync_master_yes_tx				Yes	Global	No
rpl_semi_sync_slave_enabled			Yes		Global	Yes
Rpl_semi_sync_slave_status				Yes	Global	No
rpl_semi_sync_slave_trace_level			Yes		Global	Yes
rpl_stop_slave_timeout	Yes	Yes	Yes		Global	Yes
Rsa_public_key				Yes	Global	No
safe-user-create	Yes	Yes				
schema_definition_cache	Yes	Yes	Yes		Global	Yes
secure-auth	Yes	Yes			Global	Yes
- *Variable*: secure_auth			Yes		Global	Yes
secure-file-priv	Yes	Yes			Global	No
- *Variable*: secure_file_priv			Yes		Global	No
Select_full_join				Yes	Both	No
Select_full_range_join				Yes	Both	No
Select_range				Yes	Both	No
Select_range_check				Yes	Both	No
Select_scan				Yes	Both	No
server-id	Yes	Yes			Global	Yes
- *Variable*: server_id			Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
server_uuid			Yes		Global	No
session_track_gtids	Yes	Yes	Yes		Both	Yes
session_track_schema	Yes	Yes	Yes		Both	Yes
session_track_state_change	Yes	Yes	Yes		Both	Yes
session_track_system_variables	Yes	Yes	Yes		Both	Yes
sha256_password_auto_generate_rsa_keys	Yes	Yes	Yes		Global	No
sha256_password_private_key_path			Yes		Global	No
sha256_password_proxy_users	Yes	Yes	Yes		Global	Yes
sha256_password_public_key_path			Yes		Global	No
shared_memory	Yes	Yes	Yes		Global	No
shared_memory_base_name	Yes	Yes	Yes		Global	No
show_compatibility_56	Yes	Yes	Yes		Global	Yes
show_old_temporals	Yes	Yes	Yes		Both	Yes
show-slave-auth-info	Yes	Yes				
simplified_binlog_gtid_recovery	Yes	Yes	Yes		Global	No
skip-character-set-client-handshake	Yes	Yes				
skip-concurrent-insert	Yes	Yes				
- *Variable*: concurrent_insert						
skip-event-scheduler	Yes	Yes				
skip_external_locking	Yes	Yes	Yes		Global	No
skip-grant-tables	Yes	Yes				
skip-host-cache	Yes	Yes				
skip-name-resolve	Yes	Yes			Global	No
- *Variable*: skip_name_resolve			Yes		Global	No
skip-ndbcluster	Yes	Yes				
skip-networking	Yes	Yes			Global	No
- *Variable*: skip_networking			Yes		Global	No
skip-new	Yes	Yes				
skip-partition	Yes	Yes				
skip-show-database	Yes	Yes			Global	No
- *Variable*: skip_show_database			Yes		Global	No
skip-slave-start	Yes	Yes				
skip-ssl	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
skip-stack-trace	Yes	Yes				
slave_allow_batching	Yes	Yes	Yes		Global	Yes
slave-checkpoint-group	Yes	Yes				
- *Variable*: slave_checkpoint_group						
slave_checkpoint_group	Yes	Yes	Yes		Global	Yes
slave-checkpoint-period	Yes	Yes				
- *Variable*: slave_checkpoint_period						
slave_checkpoint_period	Yes	Yes	Yes		Global	Yes
slave_compressed_protocol	Yes	Yes	Yes		Global	Yes
slave_exec_mode	Yes	Yes	Yes		Global	Yes
Slave_heartbeat_period				Yes	Global	No
Slave_last_heartbeat				Yes	Global	No
slave-load-tmpdir	Yes	Yes			Global	No
- *Variable*: slave_load_tmpdir			Yes		Global	No
slave-max-allowed-packet	Yes	Yes				
- *Variable*: slave_max_allowed_packet						
slave_max_allowed_packet			Yes		Global	Yes
slave-net-timeout	Yes	Yes			Global	Yes
- *Variable*: slave_net_timeout			Yes		Global	Yes
Slave_open_temp_tables				Yes	Global	No
slave-parallel-type	Yes	Yes				
- *Variable*: slave_parallel_type						
slave_parallel_type			Yes		Global	Yes
slave-parallel-workers	Yes	Yes				
- *Variable*: slave_parallel_workers						
slave_parallel_workers	Yes		Yes		Global	Yes
slave-pending-jobs-size-max	Yes					
- *Variable*: slave_pending_jobs_size_max						
slave_pending_jobs_size_max			Yes		Global	Yes
slave_preserve_commit_order	Yes		Yes		Global	Yes
Slave_received_heartbeats				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Slave_retried_transactions				Yes	Global	No
slave-rows-search-algorithms	Yes	Yes				
- *Variable*: slave_rows_search_algorithms						
slave_rows_search_algorithms			Yes		Global	Yes
Slave_running				Yes	Global	No
slave-skip-errors	Yes	Yes			Global	No
- *Variable*: slave_skip_errors			Yes		Global	No
slave-sql-verify-checksum	Yes	Yes				
slave_sql_verify_checksum			Yes		Global	Yes
slave_transaction_retries	Yes	Yes	Yes		Global	Yes
slave_type_conversions	Yes	Yes	Yes		Global	No
Slow_launch_threads				Yes	Both	No
slow_launch_time	Yes	Yes	Yes		Global	Yes
Slow_queries				Yes	Both	No
slow-query-log	Yes	Yes			Global	Yes
- *Variable*: slow_query_log			Yes		Global	Yes
slow_query_log_file	Yes	Yes	Yes		Global	Yes
slow-start-timeout	Yes	Yes				
socket	Yes	Yes	Yes		Global	No
sort_buffer_size	Yes	Yes	Yes		Both	Yes
Sort_merge_passes				Yes	Both	No
Sort_range				Yes	Both	No
Sort_rows				Yes	Both	No
Sort_scan				Yes	Both	No
sporadic-binlog-dump-fail	Yes	Yes				
sql_auto_is_null			Yes		Both	Yes
sql_big_selects			Yes		Both	Yes
sql_buffer_result			Yes		Both	Yes
sql_log_bin			Yes		Session	Yes
sql_log_off			Yes		Both	Yes
sql-mode	Yes	Yes			Both	Yes
- *Variable*: sql_mode			Yes		Both	Yes
sql_notes			Yes		Both	Yes
sql_quote_show_create			Yes		Both	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
sql_safe_updates			Yes		Both	Yes
sql_select_limit			Yes		Both	Yes
sql_slave_skip_counter			Yes		Global	Yes
sql_warnings			Yes		Both	Yes
ssl	Yes	Yes				
Ssl_accept_renegotiates				Yes	Global	No
Ssl_accepts				Yes	Global	No
ssl-ca	Yes	Yes			Global	No
- *Variable*: ssl_ca			Yes		Global	No
Ssl_callback_cache_hits				Yes	Global	No
ssl-capath	Yes	Yes			Global	No
- *Variable*: ssl_capath			Yes		Global	No
ssl-cert	Yes	Yes			Global	No
- *Variable*: ssl_cert			Yes		Global	No
Ssl_cipher				Yes	Both	No
ssl-cipher	Yes	Yes			Global	No
- *Variable*: ssl_cipher			Yes		Global	No
Ssl_cipher_list				Yes	Both	No
Ssl_client_connects				Yes	Global	No
Ssl_connect_renegotiates				Yes	Global	No
ssl-crl	Yes	Yes			Global	No
- *Variable*: ssl_crl			Yes		Global	No
ssl-crlpath	Yes	Yes			Global	No
- *Variable*: ssl_crlpath			Yes		Global	No
Ssl_ctx_verify_depth				Yes	Global	No
Ssl_ctx_verify_mode				Yes	Global	No
Ssl_default_timeout				Yes	Both	No
Ssl_finished_accepts				Yes	Global	No
Ssl_finished_connects				Yes	Global	No
ssl-key	Yes	Yes			Global	No
- *Variable*: ssl_key			Yes		Global	No
Ssl_server_not_after				Yes	Both	No
Ssl_server_not_before				Yes	Both	No
Ssl_session_cache_hits				Yes	Global	No
Ssl_session_cache_misses				Yes	Global	No
Ssl_session_cache_mode				Yes	Global	No
Ssl_session_cache_overflows				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Ssl_session_cache_size				Yes	Global	No
Ssl_session_cache_timeouts				Yes	Global	No
Ssl_sessions_reused				Yes	Both	No
Ssl_used_session_cache_entries				Yes	Global	No
Ssl_verify_depth				Yes	Both	No
Ssl_verify_mode				Yes	Both	No
Ssl_version				Yes	Both	No
standalone	Yes	Yes				
stored_program_cache	Yes	Yes	Yes		Global	Yes
stored_program_definition_cache	Yes	Yes	Yes		Global	Yes
super-large-pages	Yes	Yes				
super_read_only	Yes	Yes	Yes		Global	Yes
symbolic-links	Yes	Yes				
sync_binlog	Yes	Yes	Yes		Global	Yes
sync_master_info	Yes	Yes	Yes		Global	Yes
sync_relay_log	Yes	Yes	Yes		Global	Yes
sync_relay_log_info	Yes	Yes	Yes		Global	Yes
sysdate-is-now	Yes	Yes				
system_time_zone			Yes		Global	No
table_definition_cache			Yes		Global	Yes
Table_locks_immediate				Yes	Global	No
Table_locks_waited				Yes	Global	No
table_open_cache			Yes		Global	Yes
Table_open_cache_hits				Yes	Both	No
table_open_cache_instances			Yes		Global	No
Table_open_cache_misses				Yes	Both	No
Table_open_cache_overflows				Yes	Both	No
tablespace_definition_cache	Yes	Yes	Yes		Global	Yes
tc-heuristic-recover	Yes	Yes				
Tc_log_max_pages_used				Yes	Global	No
Tc_log_page_size				Yes	Global	No
Tc_log_page_waits				Yes	Global	No
temp-pool	Yes	Yes				
temptable_max_ram	Yes	Yes	Yes		Global	Yes
thread_cache_size	Yes	Yes	Yes		Global	Yes
thread_handling	Yes	Yes	Yes		Global	No
thread_stack	Yes	Yes	Yes		Global	No
Threads_cached				Yes	Global	No
Threads_connected				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Threads_created				Yes	Global	No
Threads_running				Yes	Global	No
time_format			Yes		Global	No
time_zone			Yes		Both	Yes
timestamp			Yes		Session	Yes
tls_version	Yes	Yes	Yes		Global	No
tmp_table_size	Yes	Yes	Yes		Both	Yes
tmpdir	Yes	Yes	Yes		Global	No
transaction_alloc_block_size	Yes	Yes	Yes		Both	Yes
transaction_allow_batching			Yes		Session	Yes
transaction-isolation	Yes	Yes				
- *Variable*: tx_isolation						
transaction_prealloc_size	Yes	Yes	Yes		Both	Yes
transaction-read-only	Yes	Yes				
- *Variable*: tx_read_only						
transaction_write_set_extraction	Yes		Yes		Both	Yes
tx_isolation			Yes		Both	Yes
tx_read_only			Yes		Both	Yes
unique_checks			Yes		Both	Yes
updatable_views_with_limit	Yes	Yes	Yes		Both	Yes
Uptime				Yes	Global	No
Uptime_since_flush_status				Yes	Global	No
user	Yes	Yes				
validate-password	Yes	Yes				
validate_password_check_user_name	Yes		Yes		Global	Yes
validate_password_dictionary_file			Yes		Global	Yes
validate_password_dictionary_file_last_parsed				Yes	Global	No
validate_password_dictionary_file_words_count				Yes	Global	No
validate_password_length			Yes		Global	Yes
validate_password_mixed_case_count			Yes		Global	Yes
validate_password_number_count			Yes		Global	Yes
validate_password_policy			Yes		Global	Yes
validate_password_special_char_count			Yes		Global	Yes
validate_user_plugins			Yes		Global	No
verbose	Yes	Yes				
version			Yes		Global	No
version_comment			Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
version_compile_machine			Yes		Global	No
version_compile_os			Yes		Global	No
version_tokens_session	Yes	Yes	Yes		Both	Yes
version_tokens_session_number	Yes	Yes	Yes		Both	No
wait_timeout	Yes	Yes	Yes		Both	Yes
warning_count			Yes		Session	No

[a] This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.

[b] This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.

5.1.4 Server Command Options

When you start the `mysqld` server, you can specify program options using any of the methods described in Section 4.2.3, “Specifying Program Options”. The most common methods are to provide options in an option file or on the command line. However, in most cases it is desirable to make sure that the server uses the same options each time it runs. The best way to ensure this is to list them in an option file. See Section 4.2.6, “Using Option Files”. That section also describes option file format and syntax.

`mysqld` reads options from the `[mysqld]` and `[server]` groups. `mysqld_safe` reads options from the `[mysqld]`, `[server]`, `[mysqld_safe]`, and `[safe_mysqld]` groups. `mysql.server` reads options from the `[mysqld]` and `[mysql.server]` groups.

`mysqld` accepts many command options. For a brief summary, execute this command:

```
mysqld --help
```

To see the full list, use this command:

```
mysqld --verbose --help
```

Some of the items in the list are actually system variables that can be set at server startup. These can be displayed at runtime using the `SHOW VARIABLES` statement. Some items displayed by the preceding `mysqld` command do not appear in `SHOW VARIABLES` output; this is because they are options only and not system variables.

The following list shows some of the most common server options. Additional options are described in other sections:

- Options that affect security: See Section 6.1.4, “Security-Related mysqld Options and Variables”.
- SSL-related options: See Section 6.4.5, “Command Options for Secure Connections”.
- Binary log control options: See Section 5.4.4, “The Binary Log”.
- Replication-related options: See Section 18.1.6, “Replication and Binary Logging Options and Variables”.
- Options for loading plugins such as pluggable storage engines: See Section 5.6.2, “Installing and Uninstalling Plugins”.
- Options specific to particular storage engines: See Section 15.13, “InnoDB Startup Options and System Variables” and Section 16.2.1, “MyISAM Startup Options”.

Some options control the size of buffers or caches. For a given buffer, the server might need to allocate internal data structures. These structures typically are allocated from the total memory allocated to the buffer, and the amount of space required might be platform dependent. This means that when you

assign a value to an option that controls a buffer size, the amount of space actually available might differ from the value assigned. In some cases, the amount might be less than the value assigned. It is also possible that the server will adjust a value upward. For example, if you assign a value of 0 to an option for which the minimal value is 1024, the server will set the value to 1024.

Values for buffer sizes, lengths, and stack sizes are given in bytes unless otherwise specified.

Some options take file name values. Unless otherwise specified, the default file location is the data directory if the value is a relative path name. To specify the location explicitly, use an absolute path name. Suppose that the data directory is `/var/mysql/data`. If a file-valued option is given as a relative path name, it will be located under `/var/mysql/data`. If the value is an absolute path name, its location is as given by the path name.

You can also set the values of server system variables at server startup by using variable names as options. To assign a value to a server system variable, use an option of the form `--var_name=value`. For example, `--sort_buffer_size=384M` sets the `sort_buffer_size` variable to a value of 384MB.

When you assign a value to a variable, MySQL might automatically correct the value to stay within a given range, or adjust the value to the closest permissible value if only certain values are permitted.

To restrict the maximum value to which a system variable can be set at runtime with the `SET` statement, specify this maximum by using an option of the form `--maximum-var_name=value` at server startup.

You can change the values of most system variables at runtime with the `SET` statement. See Section 13.7.4.1, “SET Syntax for Variable Assignment”.

Section 5.1.5, “Server System Variables”, provides a full description for all variables, and additional information for setting them at server startup and runtime. For information on changing system variables, see Section 5.1.1, “Configuring the Server”.

- `--help`, `-?`

Command-Line Format	`--help`

Display a short help message and exit. Use both the `--verbose` and `--help` options to see the full message.

- `--allow-suspicious-udfs`

Command-Line Format	`--allow-suspicious-udfs`	
Permitted Values	Type	boolean
	Default	`FALSE`

This option controls whether user-defined functions that have only an `xxx` symbol for the main function can be loaded. By default, the option is off and only UDFs that have at least one auxiliary symbol can be loaded; this prevents attempts at loading functions from shared object files other than those containing legitimate UDFs. See Section 28.4.2.6, “UDF Security Precautions”.

- `--ansi`

Command-Line Format	`--ansi`

Use standard (ANSI) SQL syntax instead of MySQL syntax. For more precise control over the server SQL mode, use the `--sql-mode` option instead. See Section 1.8, “MySQL Standards Compliance”, and Section 5.1.8, “Server SQL Modes”.

- `--basedir=dir_name`, `-b dir_name`

Command-Line Format	--basedir=dir_name	
System Variable	Name	basedir
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (<= 8.0.1)	Type	directory name
Permitted Values (>= 8.0.2)	Type	directory name
	Default	parent of mysqld installation directory

The path to the MySQL installation directory. This option sets the `basedir` system variable.

As of MySQL 8.0.2, the server executable determines its own full path name at startup and uses the parent of the directory in which it is located as the default `basedir` value. This in turn enables the server to use that `basedir` when searching for server-related information such as the `share` directory containing error messages.

- `--big-tables`

Command-Line Format	--big-tables	
System Variable	Name	big_tables
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Enable large result sets by saving all temporary sets in files. This option prevents most “table full” errors, but also slows down queries for which in-memory tables would suffice. The server is able to handle large result sets automatically by using memory for small temporary tables and switching to disk tables where necessary.

- `--bind-address=addr`

Command-Line Format	--bind-address=addr	
System Variable	Name	bind_address
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	*

The MySQL server listens on a single network socket for TCP/IP connections. This socket is bound to a single address, but it is possible for an address to map onto multiple network interfaces. To specify an address, use the `--bind-address=addr` option at server startup, where *`addr`* is an IPv4 or IPv6 address or a host name. If *`addr`* is a host name, the server resolves the name to an IP address and binds to that address.

The server treats different types of addresses as follows:

- If the address is `*`, the server accepts TCP/IP connections on all server host IPv6 and IPv4 interfaces if the server host supports IPv6, or accepts TCP/IP connections on all IPv4 addresses otherwise. Use this address to permit both IPv4 and IPv6 connections on all server interfaces. This value is the default.
- If the address is `0.0.0.0`, the server accepts TCP/IP connections on all server host IPv4 interfaces.
- If the address is `::`, the server accepts TCP/IP connections on all server host IPv4 and IPv6 interfaces.
- If the address is an IPv4-mapped address, the server accepts TCP/IP connections for that address, in either IPv4 or IPv6 format. For example, if the server is bound to `::ffff:127.0.0.1`, clients can connect using `--host=127.0.0.1` or `--host=::ffff:127.0.0.1`.
- If the address is a "regular" IPv4 or IPv6 address (such as `127.0.0.1` or `::1`), the server accepts TCP/IP connections only for that IPv4 or IPv6 address.

If you intend to bind the server to a specific address, be sure that the `mysql.user` grant table contains an account with administrative privileges that you can use to connect to that address. Otherwise, you will not be able to shut down the server. For example, if you bind the server to `*`, you can connect to it using all existing accounts. But if you bind the server to `::1`, it accepts connections only on that address. In that case, first make sure that the `'root'@'::1'` account is present in the `mysql.user` table so you can still connect to the server to shut it down.

- `--binlog-format={ROW|STATEMENT|MIXED}`

Command-Line Format	`--binlog-format=format`	
System Variable	**Name**	`binlog_format`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`ROW`
	Valid Values	`ROW`
		`STATEMENT`
		`MIXED`

Specify whether to use row-based, statement-based, or mixed replication. Statement-based is the default in MySQL 8.0. See Section 18.2.1, "Replication Formats".

Under some conditions, changing this variable at runtime is not possible, or causes replication to fail. See Section 5.4.4.2, "Setting The Binary Log Format", for more information.

Setting the binary logging format without enabling binary logging sets the `binlog_format` global system variable and logs a warning.

- `--character-sets-dir=dir_name`

Command-Line Format	`--character-sets-dir=dir_name`	
System Variable	**Name**	`character_sets_dir`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	directory name

The directory where character sets are installed. See Section 10.5, “Character Set Configuration”.

- `--character-set-client-handshake`

Command-Line Format	`--character-set-client-handshake`	
Permitted Values	**Type**	boolean
	Default	`TRUE`

Do not ignore character set information sent by the client. To ignore client information and use the default server character set, use `--skip-character-set-client-handshake`; this makes MySQL behave like MySQL 4.0.

- `--character-set-filesystem=`*`charset_name`*

Command-Line Format	`--character-set-filesystem=name`	
System Variable	**Name**	`character_set_filesystem`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`binary`

The file system character set. This option sets the `character_set_filesystem` system variable.

- `--character-set-server=`*`charset_name`*, `-C` *`charset_name`*

Command-Line Format	`--character-set-server`	
System Variable	**Name**	`character_set_server`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	**Type**	string
	Default	`latin1`
Permitted Values (>= 8.0.1)	**Type**	string
	Default	`utf8mb4`

Use *`charset_name`* as the default server character set. See Section 10.5, “Character Set Configuration”. If you use this option to specify a nondefault character set, you should also use `--collation-server` to specify the collation.

- `--chroot=`*`dir_name`*, `-r` *`dir_name`*

Command-Line Format	`--chroot=dir_name`	
Permitted Values	**Type**	directory name

Put the `mysqld` server in a closed environment during startup by using the `chroot()` system call. This is a recommended security measure. Use of this option somewhat limits `LOAD DATA INFILE` and `SELECT ... INTO OUTFILE`.

- `--collation-server=collation_name`

Command-Line Format	--collation-server	
System Variable	Name	collation_server
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	Type	string
	Default	latin1_swedish_ci
Permitted Values (>= 8.0.1)	Type	string
	Default	utf8mb4_0900_ai_ci

Use *`collation_name`* as the default server collation. See Section 10.5, “Character Set Configuration”.

- `--console`

Command-Line Format	--console
Platform Specific	Windows

(Windows only.) Write error log messages to `stderr` and `stdout` even if `--log-error` is specified. `mysqld` does not close the console window if this option is used.

If both `--log-error` and `--console` are specified, `--console` takes precedence. The server writes to the console, but not to the log file.

- `--core-file`

Command-Line Format	--core-file	
Permitted Values	Type	boolean
	Default	OFF

Write a core file if `mysqld` dies. The name and location of the core file is system dependent. On Linux, a core file named `core.`*`pid`* is written to the current working directory of the process, which for `mysqld` is the data directory. *`pid`* represents the process ID of the server process. On OS X, a core file named `core.`*`pid`* is written to the `/cores` directory. On Solaris, use the `coreadm` command to specify where to write the core file and how to name it.

For some systems, to get a core file you must also specify the `--core-file-size` option to `mysqld_safe`. See Section 4.3.2, “`mysqld_safe` — MySQL Server Startup Script”. On some systems, such as Solaris, you do not get a core file if you are also using the `--user` option. There might be additional restrictions or limitations. For example, it might be necessary to execute `ulimit -c unlimited` before starting the server. Consult your system documentation.

- `--daemonize`, `-D`

Command-Line Format	--daemonize[={OFF|ON}]	
Permitted Values	Type	boolean
	Default	OFF

This option causes the server to run as a traditional, forking daemon, permitting it to work with operating systems that use systemd for process control. For more information, see Section 2.5.9, “Managing MySQL Server with systemd”.

`--daemonize` is mutually exclusive with `--initialize` and `--initialize-insecure`.

As of MySQL 8.0.2, if the server is started using the `--daemonize` option and is not connected to a tty device, a default error logging option of `--log-error=""` is used in the absence of an explicit logging option, to direct error output to the default log file.

As of MySQL 8.0.2, `-D` is a synonym for `--daemonize`.

- `--datadir=dir_name`, `-h dir_name`

Command-Line Format		`--datadir=dir_name`
System Variable	Name	`datadir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The path to the MySQL server data directory. This option sets the `datadir` system variable. See the description of that variable.

- `--debug[=debug_options]`, `-# [debug_options]`

Command-Line Format		`--debug[=debug_options]`
System Variable	Name	`debug`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (Unix)	Type	string
	Default	`d:t:i:o,/tmp/mysqld.trace`
Permitted Values (Windows)	Type	string
	Default	`d:t:i:O,\mysqld.trace`

If MySQL is configured with the `-DWITH_DEBUG=1 CMake` option, you can use this option to get a trace file of what `mysqld` is doing. A typical *`debug_options`* string is `d:t:o,`*`file_name`*. The default is `d:t:i:o,/tmp/mysqld.trace` on Unix and `d:t:i:O,\mysqld.trace` on Windows.

Using `-DWITH_DEBUG=1` to configure MySQL with debugging support enables you to use the `--debug="d,parser_debug"` option when you start the server. This causes the Bison parser that is used to process SQL statements to dump a parser trace to the server's standard error output. Typically, this output is written to the error log.

This option may be given multiple times. Values that begin with `+` or `-` are added to or subtracted from the previous value. For example, `--debug=T --debug=+P` sets the value to `P:T`.

For more information, see Section 28.5.3, “The DBUG Package”.

- `--debug-sync-timeout[=N]`

Command-Line Format	--debug-sync-timeout[=#]	
Permitted Values	Type	integer

Controls whether the Debug Sync facility for testing and debugging is enabled. Use of Debug Sync requires that MySQL be configured with the -DENABLE_DEBUG_SYNC=1 CMake option (see Section 2.8.4, "MySQL Source-Configuration Options"). If Debug Sync is not compiled in, this option is not available. The option value is a timeout in seconds. The default value is 0, which disables Debug Sync. To enable it, specify a value greater than 0; this value also becomes the default timeout for individual synchronization points. If the option is given without a value, the timeout is set to 300 seconds.

For a description of the Debug Sync facility and how to use synchronization points, see MySQL Internals: Test Synchronization.

- --default-storage-engine=*type*

Command-Line Format	--default-storage-engine=name	
System Variable	Name	default_storage_engine
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	InnoDB

Set the default storage engine for tables. See Chapter 16, *Alternative Storage Engines*. This option sets the storage engine for permanent tables only. To set the storage engine for TEMPORARY tables, set the default_tmp_storage_engine system variable.

If you disable the default storage engine at server startup, you must set the default engine for both permanent and TEMPORARY tables to a different engine or the server will not start.

- --default-time-zone=*timezone*

Command-Line Format	--default-time-zone=name	
Permitted Values	Type	string

Set the default server time zone. This option sets the global time_zone system variable. If this option is not given, the default time zone is the same as the system time zone (given by the value of the system_time_zone system variable.

- --defaults-extra-file=*file_name*

Read this option file after the global option file but (on Unix) before the user option file. If the file does not exist or is otherwise inaccessible, an error occurs. *file_name* is interpreted relative to the current directory if given as a relative path name rather than a full path name. This must be the first option on the command line if it is used.

For additional information about this option, see Section 4.2.7, "Command-Line Options that Affect Option-File Handling".

- --defaults-file=*file_name*

Read only the given option file. If the file does not exist or is otherwise inaccessible, an error occurs. *file_name* is interpreted relative to the current directory if given as a relative path name rather than a full path name.

Exception: Even with `--defaults-file`, `mysqld` reads `mysqld-auto.cnf`.

Note

This must be the first option on the command line if it is used, except that if the server is started with the `--defaults-file` and `--install` (or `--install-manual`) options, `--install` (or `--install-manual`) must be first.

For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--defaults-group-suffix=`*`str`*

  Read not only the usual option groups, but also groups with the usual names and a suffix of *`str`*. For example, `mysqld` normally reads the `[mysqld]` group. If the `--defaults-group-suffix=_other` option is given, `mysqld` also reads the `[mysqld_other]` group.

  For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--delay-key-write[={OFF|ON|ALL}]`

Command-Line Format	`--delay-key-write[=name]`	
System Variable	**Name**	`delay_key_write`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`ON`
	Valid Values	`ON`
		`OFF`
		`ALL`

Specify how to use delayed key writes. Delayed key writing causes key buffers not to be flushed between writes for `MyISAM` tables. `OFF` disables delayed key writes. `ON` enables delayed key writes for those tables that were created with the `DELAY_KEY_WRITE` option. `ALL` delays key writes for all `MyISAM` tables. See Section 5.1.1, “Configuring the Server”, and Section 16.2.1, “MyISAM Startup Options”.

Note

If you set this variable to `ALL`, you should not use `MyISAM` tables from within another program (such as another MySQL server or `myisamchk`) when the tables are in use. Doing so leads to index corruption.

- `--des-key-file=`*`file_name`*

Command-Line Format	`--des-key-file=file_name`

Read the default DES keys from this file. These keys are used by the `DES_ENCRYPT()` and `DES_DECRYPT()` functions.

- `--early-plugin-load=`*`plugin_list`*

Command-Line Format	--early-plugin-load=plugin_list	
Permitted Values	Type	string
	Default	empty string

This option tells the server which plugins to load before loading mandatory built-in plugins and before storage engine initialization. If multiple `--early-plugin-load` options are given, only the last one is used.

The option value is a semicolon-separated list of *`name=plugin_library`* and *`plugin_library`* values. Each *`name`* is the name of a plugin to load, and *`plugin_library`* is the name of the library file that contains the plugin code. If a plugin library is named without any preceding plugin name, the server loads all plugins in the library. The server looks for plugin library files in the directory named by the `plugin_dir` system variable.

For example, if plugins named `myplug1` and `myplug2` have library files `myplug1.so` and `myplug2.so`, use this option to perform an early plugin load:

```
shell> mysqld --early-plugin-load="myplug1=myplug1.so;myplug2=myplug2.so"
```

Quotes are used around the argument value because otherwise a semicolon (`;`) is interpreted as a special character by some command interpreters. (Unix shells treat it as a command terminator, for example.)

Each named plugin is loaded early for a single invocation of `mysqld` only. After a restart, the plugin is not loaded early unless `--early-plugin-load` is used again.

If the server is started using `--initialize` or `--initialize-insecure`, plugins specified by `--early-plugin-load` are not loaded.

If the server is run with `--help`, plugins specified by `--early-plugin-load` are loaded but not initialized. This behavior ensures that plugin options are displayed in the help message.

The default `--early-plugin-load` value is empty. To load the `keyring_file` plugin, you must use an explicit `--early-plugin-load` option with a nonempty value.

The `InnoDB` tablespace encryption feature relies on the `keyring_file` plugin for encryption key management, and the `keyring_file` plugin must be loaded prior to storage engine initialization to facilitate `InnoDB` recovery for encrypted tables. Administrators who want the `keyring_file` plugin loaded at startup should use the appropriate nonempty option value; for example, `keyring_file.so` on Unix and Unix-like systems and `keyring_file.dll` on Windows.

For information about `InnoDB` tablespace encryption, see Section 15.7.10, “InnoDB Tablespace Encryption”. For general information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

- `--enable-named-pipe`

Command-Line Format	--enable-named-pipe
Platform Specific	Windows

Enable support for named pipes. This option applies only on Windows.

- `--event-scheduler[=value]`

Command-Line Format	--event-scheduler[=value]	
System Variable	Name	event_scheduler

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`OFF`
	Valid Values	`ON`
		`OFF`
		`DISABLED`

Enable or disable, and start or stop, the event scheduler.

For detailed information, see The `--event-scheduler` Option.

- `--exit-info[=flags]`, `-T [flags]`

Command-Line Format	`--exit-info[=flags]`	
Permitted Values	**Type**	integer

This is a bitmask of different flags that you can use for debugging the `mysqld` server. Do not use this option unless you know *exactly* what it does!

- `--external-locking`

Command-Line Format	`--external-locking`	
Permitted Values	**Type**	boolean
	Default	`FALSE`

Enable external locking (system locking), which is disabled by default. If you use this option on a system on which `lockd` does not fully work (such as Linux), it is easy for `mysqld` to deadlock.

To disable external locking explicitly, use `--skip-external-locking`.

External locking affects only `MyISAM` table access. For more information, including conditions under which it can and cannot be used, see Section 8.11.5, “External Locking”.

- `--flush`

Command-Line Format	`--flush`	
System Variable	**Name**	`flush`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Flush (synchronize) all changes to disk after each SQL statement. Normally, MySQL does a write of all changes to disk only after each SQL statement and lets the operating system handle the synchronizing to disk. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.

- `--gdb`

Command-Line Format	`--gdb`

Permitted Values	Type	boolean
	Default	FALSE

Install an interrupt handler for SIGINT (needed to stop mysqld with ^C to set breakpoints) and disable stack tracing and core file handling. See Section 28.5, “Debugging and Porting MySQL”.

- --general-log[={0|1}]

Command-Line Format	--general-log	
System Variable	Name	general_log
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Specify the initial general query log state. With no argument or an argument of 1, the --general-log option enables the log. If omitted or given with an argument of 0, the option disables the log.

- --initialize, -I

Command-Line Format	--initialize	
Permitted Values	Type	boolean
	Default	OFF

This option is used to initialize a MySQL installation by creating the data directory and populating the tables in the mysql system database. For more information, see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

When the server is started with --initialize, some functionality is unavailable that limits the statements permitted in any file named by the --init-file option. For more information, see the description of that option. In addition, the disabled_storage_engines system variable has no effect.

--initialize is mutually exclusive with --daemonize.

As of MySQL 8.0.2, -I is a synonym for --initialize.

- --initialize-insecure

Command-Line Format	--initialize-insecure	
Permitted Values	Type	boolean
	Default	OFF

This option is used to initialize a MySQL installation by creating the data directory and populating the tables in the mysql system database. This option implies --initialize. For more information, see the description of that option, and Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”.

--initialize-insecure is mutually exclusive with --daemonize.

- --init-file=*file_name*

Command-Line Format	--init-file=file_name

System Variable	Name	init_file
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

Read SQL statements from this file at startup. Each statement must be on a single line and should not include comments.

If the server is started with the `--initialize` or `--initialize-insecure` option, it operates in bootstap mode and some functionality is unavailable that limits the statements permitted in the file. These include statements that relate to account management (such as `CREATE USER` or `GRANT`), replication, and global transaction identifiers. See Section 18.1.3, “Replication with Global Transaction Identifiers”.

- `--innodb-xxx`

  Set an option for the `InnoDB` storage engine. The `InnoDB` options are listed in Section 15.13, “InnoDB Startup Options and System Variables”.

- `--install [service_name]`

Command-Line Format	--install [service_name]
Platform Specific	Windows

(Windows only) Install the server as a Windows service that starts automatically during Windows startup. The default service name is `MySQL` if no *`service_name`* value is given. For more information, see Section 2.3.5.8, “Starting MySQL as a Windows Service”.

Note

If the server is started with the `--defaults-file` and `--install` options, `--install` must be first.

- `--install-manual [service_name]`

Command-Line Format	--install-manual [service_name]
Platform Specific	Windows

(Windows only) Install the server as a Windows service that must be started manually. It does not start automatically during Windows startup. The default service name is `MySQL` if no *`service_name`* value is given. For more information, see Section 2.3.5.8, “Starting MySQL as a Windows Service”.

Note

If the server is started with the `--defaults-file` and `--install-manual` options, `--install-manual` must be first.

- `--language=lang_name, -L lang_name`

Deprecated	5.6.1, by lc-messages-dir
Command-Line Format	--language=name
System Variable	Name language
	Variable Scope Global

	Dynamic Variable	No
Permitted Values	**Type**	directory name
	Default	`/usr/local/mysql/share/mysql/english/`

The language to use for error messages. *`lang_name`* can be given as the language name or as the full path name to the directory where the language files are installed. See Section 10.2, “Setting the Error Message Language”.

`--lc-messages-dir` and `--lc-messages` should be used rather than `--language`, which is deprecated (and handled as an alias for `--lc-messages-dir`). The `--language` option will be removed in a future MySQL release.

- `--large-pages`

Command-Line Format	`--large-pages`	
System Variable	**Name**	`large_pages`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Linux	
Permitted Values (Linux)	**Type**	boolean
	Default	`FALSE`

Some hardware/operating system architectures support memory pages greater than the default (usually 4KB). The actual implementation of this support depends on the underlying hardware and operating system. Applications that perform a lot of memory accesses may obtain performance improvements by using large pages due to reduced Translation Lookaside Buffer (TLB) misses.

MySQL supports the Linux implementation of large page support (which is called HugeTLB in Linux). See Section 8.12.3.2, “Enabling Large Page Support”. For Solaris support of large pages, see the description of the `--super-large-pages` option.

`--large-pages` is disabled by default.

- `--lc-messages=`*`locale_name`*

Command-Line Format	`--lc-messages=name`	
System Variable	**Name**	`lc_messages`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`en_US`

The locale to use for error messages. The default is `en_US`. The server converts the argument to a language name and combines it with the value of `--lc-messages-dir` to produce the location for the error message file. See Section 10.2, “Setting the Error Message Language”.

- `--lc-messages-dir=`*`dir_name`*

Command-Line Format	`--lc-messages-dir=dir_name`

System Variable	Name	lc_messages_dir
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The directory where error messages are located. The server uses the value together with the value of `--lc-messages` to produce the location for the error message file. See Section 10.2, “Setting the Error Message Language”.

- `--local-service`

Command-Line Format	--local-service

(Windows only) A `--local-service` option following the service name causes the server to run using the `LocalService` Windows account that has limited system privileges. If both `--defaults-file` and `--local-service` are given following the service name, they can be in any order. See Section 2.3.5.8, “Starting MySQL as a Windows Service”.

- `--log-error[=file_name]`

Command-Line Format	--log-error[=file_name]	
System Variable	Name	log_error
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

Log errors and startup messages to this file. See Section 5.4.2, “The Error Log”. If the file name has no extension, the server adds an extension of `.err`. If you omit the file name, the default log file on Unix and Unix-like systems is *`host_name`*`.err` in the data directory. The default on Windows is *`host_name`*`.err` in the data directory, unless the `--pid-file` option is specified. In that case, the default name is the PID file base name with a suffix of `.err` in the data directory.

As of MySQL 8.0.2, if error log output cannot be redirected to the `--log-error` option value, an error occurs and startup fails.

- `--log-isam[=file_name]`

Command-Line Format	--log-isam[=file_name]	
Permitted Values	Type	file name

Log all `MyISAM` changes to this file (used only when debugging `MyISAM`).

- `--log-output=value,...`

Command-Line Format	--log-output=name	
System Variable	Name	log_output
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	set
	Default	`FILE`
	Valid Values	`TABLE`
		`FILE`
		`NONE`

This option determines the destination for general query log and slow query log output. The option value can be given as one or more of the words `TABLE`, `FILE`, or `NONE`. `TABLE` select logging to the `general_log` and `slow_log` tables in the `mysql` database as a destination. `FILE` selects logging to log files as a destination. `NONE` disables logging. If `NONE` is present in the option value, it takes precedence over any other words that are present. `TABLE` and `FILE` can both be given to select to both log output destinations.

This option selects log output destinations, but does not enable log output. To do that, use the `--general_log` and `--slow_query_log` options. For `FILE` logging, the `--general_log_file` and `-slow_query_log_file` options determine the log file location. For more information, see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”.

- `--log-queries-not-using-indexes`

Command-Line Format	`--log-queries-not-using-indexes`	
System Variable	**Name**	`log_queries_not_using_indexes`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

If you are using this option with the slow query log enabled, queries that are expected to retrieve all rows are logged. See Section 5.4.5, “The Slow Query Log”. This option does not necessarily mean that no index is used. For example, a query that uses a full index scan uses an index but would be logged because the index would not limit the number of rows.

- `--log-raw`

Command-Line Format	`--log-raw[=value]`	
Permitted Values	**Type**	boolean
	Default	`OFF`

Passwords in certain statements written to the general query log, slow query log, and binary log are rewritten by the server not to occur literally in plain text. Password rewriting can be suppressed for the general query log by starting the server with the `--log-raw` option. This option may be useful for diagnostic purposes, to see the exact text of statements as received by the server, but for security reasons is not recommended for production use.

If a query rewrite plugin is installed, the `--log-raw` option affects statement logging as follows:

- Without `--log-raw`, the server logs the statement returned by the query rewrite plugin. This may differ from the statement as received.
- With `--log-raw`, the server logs the original statement as received.

For more information, see Section 6.1.2.3, “Passwords and Logging”.

- `--log-short-format`

Command-Line Format	`--log-short-format`	
Permitted Values	Type	boolean
	Default	`FALSE`

Log less information to the slow query log, if it has been activated.

- `--log-tc=file_name`

Command-Line Format	`--log-tc=file_name`	
Permitted Values	Type	file name
	Default	`tc.log`

The name of the memory-mapped transaction coordinator log file (for XA transactions that affect multiple storage engines when the binary log is disabled). The default name is `tc.log`. The file is created under the data directory if not given as a full path name. This option is unused.

- `--log-tc-size=size`

Command-Line Format	`--log-tc-size=#`	
Permitted Values (32-bit platforms)	Type	integer
	Default	`6 * page size`
	Min Value	`6 * page size`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	Type	integer
	Default	`6 * page size`
	Min Value	`6 * page size`
	Max Value	`18446744073709551615`

The size in bytes of the memory-mapped transaction coordinator log. The default and minimum values are 6 times the page size, and the value must be a multiple of the page size.

- `--log-warnings[=level]`, `-W [level]`

Deprecated	5.7.2	
Command-Line Format	`--log-warnings[=#]`	
System Variable	Name	`log_warnings`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	`2`
	Min Value	`0`

	Max Value	4294967295
Permitted Values (64-bit platforms)	**Type**	integer
	Default	2
	Min Value	0
	Max Value	18446744073709551615

Note

The `log_error_verbosity` system variable is preferred over, and should be used instead of, the `--log-warnings` option or `log_warnings` system variable. For more information, see the descriptions of `log_error_verbosity` and `log_warnings`. The `--log-warnings` command-line option and `log_warnings` system variable are deprecated and will be removed in a future MySQL release.

Whether to produce additional warning messages to the error log. This option is enabled by default with a value of 2. To disable it, use `--log-warnings=0`. Specifying the option without a *`level`* value increments the current value by 1. The server logs messages about statements that are unsafe for statement-based logging if the value is greater than 0. Aborted connections and access-denied errors for new connection attempts are logged if the value is greater than 1. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

- `--low-priority-updates`

Command-Line Format	`--low-priority-updates`	
System Variable	**Name**	`low_priority_updates`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`FALSE`

Give table-modifying operations (`INSERT`, `REPLACE`, `DELETE`, `UPDATE`) lower priority than selects. This can also be done using `{INSERT | REPLACE | DELETE | UPDATE} LOW_PRIORITY ...` to lower the priority of only one query, or by `SET LOW_PRIORITY_UPDATES=1` to change the priority in one thread. This affects only storage engines that use only table-level locking (`MyISAM`, `MEMORY`, `MERGE`). See Section 8.11.2, “Table Locking Issues”.

- `--min-examined-row-limit=`*`number`*

Command-Line Format	`--min-examined-row-limit=#`	
System Variable	**Name**	`min_examined_row_limit`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	0

	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709551615`

When this option is set, queries which examine fewer than *`number`* rows are not written to the slow query log. The default is 0.

- `--memlock`

Command-Line Format	`--memlock`	
Permitted Values	**Type**	boolean
	Default	`FALSE`

Lock the `mysqld` process in memory. This option might help if you have a problem where the operating system is causing `mysqld` to swap to disk.

`--memlock` works on systems that support the `mlockall()` system call; this includes Solaris, most Linux distributions that use a 2.4 or higher kernel, and perhaps other Unix systems. On Linux systems, you can tell whether or not `mlockall()` (and thus this option) is supported by checking to see whether or not it is defined in the system `mman.h` file, like this:

```
shell> grep mlockall /usr/include/sys/mman.h
```

If `mlockall()` is supported, you should see in the output of the previous command something like the following:

```
extern int mlockall (int __flags) __THROW;
```

Important

Use of this option may require you to run the server as `root`, which, for reasons of security, is normally not a good idea. See Section 6.1.5, “How to Run MySQL as a Normal User”.

On Linux and perhaps other systems, you can avoid the need to run the server as `root` by changing the `limits.conf` file. See the notes regarding the memlock limit in Section 8.12.3.2, “Enabling Large Page Support”.

You must not try to use this option on a system that does not support the `mlockall()` system call; if you do so, `mysqld` will very likely crash as soon as you try to start it.

- `--myisam-block-size=N`

Command-Line Format	`--myisam-block-size=#`	
Permitted Values	**Type**	integer
	Default	`1024`

	Min Value	1024
	Max Value	16384

The block size to be used for MyISAM index pages.

- `--myisam-recover-options[=option[,option]...]]`

Command-Line Format	`--myisam-recover-options[=name]`	
Permitted Values	**Type**	enumeration
	Default	OFF
	Valid Values	OFF
		DEFAULT
		BACKUP
		FORCE
		QUICK

Set the MyISAM storage engine recovery mode. The option value is any combination of the values of OFF, DEFAULT, BACKUP, FORCE, or QUICK. If you specify multiple values, separate them by commas. Specifying the option with no argument is the same as specifying DEFAULT, and specifying with an explicit value of "" disables recovery (same as a value of OFF). If recovery is enabled, each time mysqld opens a MyISAM table, it checks whether the table is marked as crashed or was not closed properly. (The last option works only if you are running with external locking disabled.) If this is the case, mysqld runs a check on the table. If the table was corrupted, mysqld attempts to repair it.

The following options affect how the repair works.

Option	Description
OFF	No recovery.
DEFAULT	Recovery without backup, forcing, or quick checking.
BACKUP	If the data file was changed during recovery, save a backup of the *tbl_name*.MYD file as *tbl_name-datetime*.BAK.
FORCE	Run recovery even if we would lose more than one row from the .MYD file.
QUICK	Do not check the rows in the table if there are not any delete blocks.

Before the server automatically repairs a table, it writes a note about the repair to the error log. If you want to be able to recover from most problems without user intervention, you should use the options BACKUP,FORCE. This forces a repair of a table even if some rows would be deleted, but it keeps the old data file as a backup so that you can later examine what happened.

See Section 16.2.1, “MyISAM Startup Options”.

- `--no-defaults`

Do not read any option files. If program startup fails due to reading unknown options from an option file, `--no-defaults` can be used to prevent them from being read. This must be the first option on the command line if it is used.

For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--old-alter-table`

Command-Line Format	--old-alter-table	
System Variable	Name	old_alter_table
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

When this option is given, the server does not use the optimized method of processing an ALTER TABLE operation. It reverts to using a temporary table, copying over the data, and then renaming the temporary table to the original, as used by MySQL 5.0 and earlier. For more information on the operation of ALTER TABLE, see Section 13.1.7, “ALTER TABLE Syntax”.

- --old-style-user-limits

Command-Line Format	--old-style-user-limits	
Permitted Values	Type	boolean
	Default	FALSE

Enable old-style user limits. (Before MySQL 5.0.3, account resource limits were counted separately for each host from which a user connected rather than per account row in the user table.) See Section 6.3.6, “Setting Account Resource Limits”.

- --open-files-limit=*count*

Command-Line Format	--open-files-limit=#	
System Variable	Name	open_files_limit
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	5000, with possible adjustment
	Min Value	0
	Max Value	platform dependent

Changes the number of file descriptors available to mysqld. You should try increasing the value of this option if mysqld gives you the error Too many open files. mysqld uses the option value to reserve descriptors with setrlimit(). Internally, the maximum value for this option is the maximum unsigned integer value, but the actual maximum is platform dependent. If the requested number of file descriptors cannot be allocated, mysqld writes a warning to the error log.

mysqld may attempt to allocate more than the requested number of descriptors (if they are available), using the values of max_connections and table_open_cache to estimate whether more descriptors will be needed.

On Unix, the value cannot be set greater than ulimit -n.

- --performance-schema-xxx

Configure a Performance Schema option. For details, see Section 25.13, “Performance Schema Command Options”.

- `--pid-file=file_name`

Command-Line Format	--pid-file=file_name	
System Variable	Name	pid_file
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The path name of the process ID file. The server creates the file in the data directory unless an absolute path name is given to specify a different directory. This file is used by other programs such as `mysqld_safe` to determine the server's process ID.

- `--plugin-xxx`

Specifies an option that pertains to a server plugin. For example, many storage engines can be built as plugins, and for such engines, options for them can be specified with a `--plugin` prefix. Thus, the `--innodb_file_per_table` option for `InnoDB` can be specified as `--plugin-innodb_file_per_table`.

For boolean options that can be enabled or disabled, the `--skip` prefix and other alternative formats are supported as well (see Section 4.2.5, “Program Option Modifiers”). For example, `--skip-plugin-innodb_file_per_table` disables `innodb_file_per_table`.

The rationale for the `--plugin` prefix is that it enables plugin options to be specified unambiguously if there is a name conflict with a built-in server option. For example, were a plugin writer to name a plugin “sql” and implement a “mode” option, the option name might be `--sql-mode`, which would conflict with the built-in option of the same name. In such cases, references to the conflicting name are resolved in favor of the built-in option. To avoid the ambiguity, users can specify the plugin option as `--plugin-sql-mode`. Use of the `--plugin` prefix for plugin options is recommended to avoid any question of ambiguity.

- `--plugin-load=plugin_list`

Command-Line Format	--plugin-load=plugin_list	
Permitted Values	Type	string

This option tells the server to load the named plugins at startup. If multiple `--plugin-load` options are given, only the last one is used. Additional plugins to load may be specified using `--plugin-load-add` options.

The option value is a semicolon-separated list of *`name=plugin_library`* and *`plugin_library`* values. Each *`name`* is the name of a plugin to load, and *`plugin_library`* is the name of the library file that contains the plugin code. If a plugin library is named without any preceding plugin name, the server loads all plugins in the library. The server looks for plugin library files in the directory named by the `plugin_dir` system variable.

For example, if plugins named `myplug1` and `myplug2` have library files `myplug1.so` and `myplug2.so`, use this option to perform an early plugin load:

```
shell> mysqld --plugin-load="myplug1=myplug1.so;myplug2=myplug2.so"
```

Quotes are used around the argument value here because otherwise semicolon (`;`) is interpreted as a special character by some command interpreters. (Unix shells treat it as a command terminator, for example.)

Each named plugin is loaded for a single invocation of `mysqld` only. After a restart, the plugin is not loaded unless `--plugin-load` is used again. This is in contrast to `INSTALL PLUGIN`, which adds an entry to the `mysql.plugins` table to cause the plugin to be loaded for every normal server startup.

Under normal startup, the server determines which plugins to load by reading the `mysql.plugins` system table. If the server is started with the `--skip-grant-tables` option, it does not consult the `mysql.plugins` table and does not load plugins listed there. `--plugin-load` enables plugins to be loaded even when `--skip-grant-tables` is given. `--plugin-load` also enables plugins to be loaded at startup that cannot be loaded at runtime.

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

- `--plugin-load-add=`*`plugin_list`*

Command-Line Format	`--plugin-load-add=plugin_list`	
Permitted Values	**Type**	string

This option complements the `--plugin-load` option. `--plugin-load-add` adds a plugin or plugins to the set of plugins to be loaded at startup. The argument format is the same as for `--plugin-load`. `--plugin-load-add` can be used to avoid specifying a large set of plugins as a single long unwieldy `--plugin-load` argument.

`--plugin-load-add` can be given in the absence of `--plugin-load`, but any instance of `--plugin-load-add` that appears before `--plugin-load`. has no effect because `--plugin-load` resets the set of plugins to load. In other words, these options:

```
--plugin-load=x --plugin-load-add=y
```

are equivalent to this option:

```
--plugin-load="x;y"
```

But these options:

```
--plugin-load-add=y --plugin-load=x
```

are equivalent to this option:

```
--plugin-load=x
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

- `--port=`*`port_num`*, `-P` *`port_num`*

Command-Line Format	`--port=#`	
System Variable	**Name**	`port`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`3306`
	Min Value	`0`
	Max Value	`65535`

The port number to use when listening for TCP/IP connections. On Unix and Unix-like systems, the port number must be 1024 or higher unless the server is started by the `root` system user.

- `--port-open-timeout=`*`num`*

Command-Line Format	`--port-open-timeout=#`	
Permitted Values	**Type**	integer
	Default	`0`

On some systems, when the server is stopped, the TCP/IP port might not become available immediately. If the server is restarted quickly afterward, its attempt to reopen the port can fail. This option indicates how many seconds the server should wait for the TCP/IP port to become free if it cannot be opened. The default is not to wait.

- `--print-defaults`

Print the program name and all options that it gets from option files. Password values are masked. This must be the first option on the command line if it is used, except that it may be used immediately after `--defaults-file` or `--defaults-extra-file`.

For additional information about this option, see Section 4.2.7, “Command-Line Options that Affect Option-File Handling”.

- `--remove [`*`service_name`*`]`

Command-Line Format	`--remove [service_name]`
Platform Specific	Windows

(Windows only) Remove a MySQL Windows service. The default service name is `MySQL` if no *`service_name`* value is given. For more information, see Section 2.3.5.8, “Starting MySQL as a Windows Service”.

- `--safe-user-create`

Command-Line Format	`--safe-user-create`	
Permitted Values	**Type**	boolean
	Default	`FALSE`

If this option is enabled, a user cannot create new MySQL users by using the `GRANT` statement unless the user has the `INSERT` privilege for the `mysql.user` table or any column in the table. If you want a user to have the ability to create new users that have those privileges that the user has the right to grant, you should grant the user the following privilege:

```
GRANT INSERT(user) ON mysql.user TO 'user_name'@'host_name';
```

This ensures that the user cannot change any privilege columns directly, but has to use the `GRANT` statement to give privileges to other users.

- `--secure-auth`

Deprecated	5.7.5	
Command-Line Format	`--secure-auth`	
System Variable	**Name**	`secure_auth`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`
	Valid Values	`ON`

This option causes the server to block connections by clients that attempt to use accounts that have passwords stored in the old (pre-4.1) format. Use it to prevent all use of passwords employing the old format (and hence insecure communication over the network).

This option is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it (`--skip-secure-auth`, `--secure-auth=0`) produces an error.

Server startup fails with an error if this option is enabled and the privilege tables are in pre-4.1 format.

The `mysql` client also has a `--secure-auth` option, which prevents connections to a server if the server requires a password in old format for the client account.

Note

Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- `--secure-file-priv=`*`dir_name`*

Command-Line Format	`--secure-file-priv=dir_name`	
System Variable	**Name**	`secure_file_priv`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`platform specific`
	Valid Values	`empty`
		`dirname`
		`NULL`

This option sets the `secure_file_priv` system variable, which is used to limit the effect of data import and export operations, such as those performed by the `LOAD DATA` and `SELECT ... INTO`

`OUTFILE` statements and the `LOAD_FILE()` function. For more information, see the description of `secure_file_priv`.

- `--shared-memory`

Command-Line Format	`--shared-memory[={0,1}]`	
System Variable	Name	`shared_memory`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Windows	
Permitted Values	Type	boolean
	Default	`FALSE`

Enable shared-memory connections by local clients. This option is available only on Windows.

- `--shared-memory-base-name=`*name*

Command-Line Format	`--shared-memory-base-name=name`	
System Variable	Name	`shared_memory_base_name`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Windows	
Permitted Values	Type	string
	Default	`MYSQL`

The name of shared memory to use for shared-memory connections. This option is available only on Windows. The default name is `MYSQL`. The name is case sensitive.

- `--skip-concurrent-insert`

Turn off the ability to select and insert at the same time on `MyISAM` tables. (This is to be used only if you think you have found a bug in this feature.) See Section 8.11.3, “Concurrent Inserts”.

- `--skip-event-scheduler`

Command-Line Format	`--skip-event-scheduler`
	`--disable-event-scheduler`

Turns the Event Scheduler `OFF`. This is not the same as disabling the Event Scheduler, which requires setting `--event-scheduler=DISABLED`; see The `--event-scheduler` Option, for more information.

- `--skip-grant-tables`

This option causes the server to start without using the privilege system at all, which gives anyone with access to the server *unrestricted access to all databases*. You can cause a running server to start using the grant tables again by executing `mysqladmin flush-privileges` or `mysqladmin reload` command from a system shell, or by issuing a MySQL `FLUSH PRIVILEGES` statement after connecting to the server.

This option also causes the server to suppress during its startup sequence the loading of user-defined functions (UDFs), scheduled events, and plugins that were installed with the `INSTALL PLUGIN` statement. To cause plugins to be loaded anyway, use the `--plugin-load` option. `--skip-grant-tables` also causes the `disabled_storage_engines` system variable to have no effect.

This option does not cause loading of server components to be suppressed during server startup.

`FLUSH PRIVILEGES` might be executed implicitly by other actions performed after startup. For example, `mysql_upgrade` flushes the privileges during the upgrade procedure.

- `--skip-host-cache`

  Disable use of the internal host cache for faster name-to-IP resolution. In this case, the server performs a DNS lookup every time a client connects. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.

  Use of `--skip-host-cache` is similar to setting the `host_cache_size` system variable to 0, but `host_cache_size` is more flexible because it can also be used to resize, enable, or disable the host cache at runtime, not just at server startup.

  If you start the server with `--skip-host-cache`, that does not prevent changes to the value of `host_cache_size`, but such changes have no effect and the cache is not re-enabled even if `host_cache_size` is set larger than 0.

- `--skip-innodb`

  Disable the `InnoDB` storage engine. In this case, because the default storage engine is `InnoDB`, the server will not start unless you also use `--default-storage-engine` and `--default-tmp-storage-engine` to set the default to some other engine for both permanent and `TEMPORARY` tables.

  The `InnoDB` storage engine cannot be disabled, and the `--skip-innodb` option is deprecated and has no effect. Its use results in a warning. This option will be removed in a future MySQL release.

- `--skip-name-resolve`

  Do not resolve host names when checking client connections. Use only IP addresses. If you use this option, all `Host` column values in the grant tables must be IP addresses. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.

  Depending on the network configuration of your system and the `Host` values for your accounts, clients may need to connect using an explicit `--host` option, such as `--host=127.0.0.1` or `--host=::1`.

  An attempt to connect to the host `127.0.0.1` normally resolves to the `localhost` account. However, this fails if the server is run with the `--skip-name-resolve` option, so make sure that an account exists that can accept a connection. For example, to be able to connect as `root` using `--host=127.0.0.1` or `--host=::1`, create these accounts:

  ```
  CREATE USER 'root'@'127.0.0.1' IDENTIFIED BY 'root-password';
  CREATE USER 'root'@'::1' IDENTIFIED BY 'root-password';
  ```

- `--skip-networking`

  Do not listen for TCP/IP connections at all. All interaction with `mysqld` must be made using named pipes or shared memory (on Windows) or Unix socket files (on Unix). This option is highly recommended for systems where only local clients are permitted. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.

- `--ssl*`

  Options that begin with `--ssl` specify whether to permit clients to connect using SSL and indicate where to find SSL keys and certificates. See Section 6.4.5, “Command Options for Secure Connections”.

- `--standalone`

Command-Line Format	`--standalone`	
Platform Specific	Windows	

  Available on Windows only; instructs the MySQL server not to run as a service.

- `--super-large-pages`

Command-Line Format	`--super-large-pages`	
Platform Specific	Solaris	
Permitted Values (Solaris)	**Type**	boolean
	Default	`FALSE`

  Standard use of large pages in MySQL attempts to use the largest size supported, up to 4MB. Under Solaris, a “super large pages” feature enables uses of pages up to 256MB. This feature is available for recent SPARC platforms. It can be enabled or disabled by using the `--super-large-pages` or `--skip-super-large-pages` option.

- `--symbolic-links`, `--skip-symbolic-links`

Deprecated	8.0.2	
Command-Line Format	`--symbolic-links`	
Permitted Values (<= 8.0.1)	**Type**	boolean
	Default	`ON`
Permitted Values (>= 8.0.2)	**Type**	boolean
	Default	`OFF`

  Enable or disable symbolic link support. On Unix, enabling symbolic links means that you can link a `MyISAM` index file or data file to another directory with the `INDEX DIRECTORY` or `DATA DIRECTORY` options of the `CREATE TABLE` statement. If you delete or rename the table, the files that its symbolic links point to also are deleted or renamed. See Section 8.12.2.2, “Using Symbolic Links for MyISAM Tables on Unix”.

  Note

  As of MySQL 8.0.2, symbolic link support, along with the the `--symbolic-links` option that controls it, is deprecated and will be removed in a future version of MySQL. In addition, the option is disabled by default. The related `have_symlink` system variable also is deprecated and will be removed in a future version of MySQL.

  This option has no meaning on Windows.

- `--skip-show-database`

Command-Line Format	`--skip-show-database`	
System Variable	**Name**	`skip_show_database`

	Variable Scope	Global
	Dynamic Variable	No

This option sets the `skip_show_database` system variable that controls who is permitted to use the `SHOW DATABASES` statement. See Section 5.1.5, “Server System Variables”.

- `--skip-stack-trace`

Command-Line Format	`--skip-stack-trace`

Do not write stack traces. This option is useful when you are running `mysqld` under a debugger. On some systems, you also must use this option to get a core file. See Section 28.5, “Debugging and Porting MySQL”.

- `--slow-query-log[={0|1}]`

Command-Line Format	`--slow-query-log`	
System Variable	**Name**	`slow_query_log`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Specify the initial slow query log state. With no argument or an argument of 1, the `--slow-query-log` option enables the log. If omitted or given with an argument of 0, the option disables the log.

- `--slow-start-timeout=`*`timeout`*

Command-Line Format	`--slow-start-timeout=#`	
Permitted Values (Windows)	**Type**	integer
	Default	`15000`

This option controls the Windows service control manager's service start timeout. The value is the maximum number of milliseconds that the service control manager waits before trying to kill the windows service during startup. The default value is 15000 (15 seconds). If the MySQL service takes too long to start, you may need to increase this value. A value of 0 means there is no timeout.

- `--socket=`*`path`*

Command-Line Format	`--socket={file_name|pipe_name}`	
System Variable	**Name**	`socket`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`/tmp/mysql.sock`

On Unix, this option specifies the Unix socket file to use when listening for local connections. The default value is `/tmp/mysql.sock`. If this option is given, the server creates the file in the data

directory unless an absolute path name is given to specify a different directory. On Windows, the option specifies the pipe name to use when listening for local connections that use a named pipe. The default value is `MySQL` (not case sensitive).

- `--sql-mode=value[,value[,value...]]`

Command-Line Format	`--sql-mode=name`	
System Variable	**Name**	`sql_mode`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	set
	Default	`ONLY_FULL_GROUP_BY STRICT_TRANS_TABLES NO_ZERO_IN_DATE NO_ZERO_DATE ERROR_FOR_DIVISION_BY_ZERO NO_AUTO_CREATE_USER NO_ENGINE_SUBSTITUTION`
	Valid Values	`ALLOW_INVALID_DATES`
		`ANSI_QUOTES`
		`ERROR_FOR_DIVISION_BY_ZERO`
		`HIGH_NOT_PRECEDENCE`
		`IGNORE_SPACE`
		`NO_AUTO_CREATE_USER`
		`NO_AUTO_VALUE_ON_ZERO`
		`NO_BACKSLASH_ESCAPES`
		`NO_DIR_IN_CREATE`
		`NO_ENGINE_SUBSTITUTION`
		`NO_FIELD_OPTIONS`
		`NO_KEY_OPTIONS`
		`NO_TABLE_OPTIONS`
		`NO_UNSIGNED_SUBTRACTION`
		`NO_ZERO_DATE`
		`NO_ZERO_IN_DATE`
		`ONLY_FULL_GROUP_BY`
		`PAD_CHAR_TO_FULL_LENGTH`
		`PIPES_AS_CONCAT`
		`REAL_AS_FLOAT`
		`STRICT_ALL_TABLES`
		`STRICT_TRANS_TABLES`
Permitted Values (>= 8.0.1)	**Type**	set
	Default	`ONLY_FULL_GROUP_BY STRICT_TRANS_TABLES NO_ZERO_IN_DATE NO_ZERO_DATE ERROR_FOR_DIVISION_BY_ZERO NO_AUTO_CREATE_USER NO_ENGINE_SUBSTITUTION`
	Valid Values	`ALLOW_INVALID_DATES`
		`ANSI_QUOTES`

		ERROR_FOR_DIVISION_BY_ZERO
		HIGH_NOT_PRECEDENCE
		IGNORE_SPACE
		NO_AUTO_CREATE_USER
		NO_AUTO_VALUE_ON_ZERO
		NO_BACKSLASH_ESCAPES
		NO_DIR_IN_CREATE
		NO_ENGINE_SUBSTITUTION
		NO_FIELD_OPTIONS
		NO_KEY_OPTIONS
		NO_TABLE_OPTIONS
		NO_UNSIGNED_SUBTRACTION
		NO_ZERO_DATE
		NO_ZERO_IN_DATE
		ONLY_FULL_GROUP_BY
		PAD_CHAR_TO_FULL_LENGTH
		PIPES_AS_CONCAT
		REAL_AS_FLOAT
		STRICT_ALL_TABLES
		STRICT_TRANS_TABLES
		TIME_TRUNCATE_FRACTIONAL

Set the SQL mode. See Section 5.1.8, “Server SQL Modes”.

Note

MySQL installation programs may configure the SQL mode during the installation process.

If the SQL mode differs from the default or from what you expect, check for a setting in an option file that the server reads at startup.

- `--sysdate-is-now`

Command-Line Format	`--sysdate-is-now`	
Permitted Values	**Type**	boolean
	Default	`FALSE`

`SYSDATE()` by default returns the time at which it executes, not the time at which the statement in which it occurs begins executing. This differs from the behavior of `NOW()`. This option causes `SYSDATE()` to be an alias for `NOW()`. For information about the implications for binary logging and replication, see the description for `SYSDATE()` in Section 12.7, “Date and Time Functions” and for `SET TIMESTAMP` in Section 5.1.5, “Server System Variables”.

- `--tc-heuristic-recover={COMMIT|ROLLBACK}`

Command-Line Format	`--tc-heuristic-recover=name`	
Permitted Values	**Type**	enumeration
	Default	`COMMIT`

	Valid Values	`COMMIT`
		`ROLLBACK`

The type of decision to use in the heuristic recovery process. To use this option, two or more storage engines that support XA transactions must be installed.

- `--temp-pool`

Deprecated	5.7.18	
Removed	8.0.1	
Command-Line Format	`--temp-pool`	
Permitted Values	**Type**	boolean
	Default	`TRUE`

This option is obsolete and was removed in MySQL 8.0.1.

- `--transaction-isolation=`*`level`*

Command-Line Format	`--transaction-isolation=name`	
Permitted Values	**Type**	enumeration
	Default	`REPEATABLE-READ`
	Valid Values	`READ-UNCOMMITTED`
		`READ-COMMITTED`
		`REPEATABLE-READ`
		`SERIALIZABLE`

Sets the default transaction isolation level. The `level` value can be `READ-UNCOMMITTED`, `READ-COMMITTED`, `REPEATABLE-READ`, or `SERIALIZABLE`. See Section 13.3.6, “SET TRANSACTION Syntax”.

The default transaction isolation level can also be set at runtime using the `SET TRANSACTION` statement or by setting the `tx_isolation` system variable.

- `--transaction-read-only`

Command-Line Format	`--transaction-read-only`	
Permitted Values	**Type**	boolean
	Default	`OFF`

Sets the default transaction access mode. By default, read-only mode is disabled, so the mode is read/write.

To set the default transaction access mode at runtime, use the `SET TRANSACTION` statement or set the `tx_read_only` system variable. See Section 13.3.6, “SET TRANSACTION Syntax”.

- `--tmpdir=`*`dir_name`*, `-t` *`dir_name`*

Command-Line Format	`--tmpdir=dir_name`	
System Variable	**Name**	`tmpdir`
	Variable Scope	Global
	Dynamic Variable	No

Permitted Values	Type	directory name

The path of the directory to use for creating temporary files. It might be useful if your default `/tmp` directory resides on a partition that is too small to hold temporary tables. This option accepts several paths that are used in round-robin fashion. Paths should be separated by colon characters (`:`) on Unix and semicolon characters (`;`) on Windows. If the MySQL server is acting as a replication slave, you should not set `--tmpdir` to point to a directory on a memory-based file system or to a directory that is cleared when the server host restarts. For more information about the storage location of temporary files, see Section B.5.3.5, "Where MySQL Stores Temporary Files". A replication slave needs some of its temporary files to survive a machine restart so that it can replicate temporary tables or `LOAD DATA INFILE` operations. If files in the temporary file directory are lost when the server restarts, replication fails.

- `--user={user_name|user_id}`, `-u {user_name|user_id}`

Command-Line Format	`--user=name`	
Permitted Values	Type	string

Run the `mysqld` server as the user having the name *`user_name`* or the numeric user ID *`user_id`*. ("User" in this context refers to a system login account, not a MySQL user listed in the grant tables.)

This option is *mandatory* when starting `mysqld` as `root`. The server changes its user ID during its startup sequence, causing it to run as that particular user rather than as `root`. See Section 6.1.1, "Security Guidelines".

To avoid a possible security hole where a user adds a `--user=root` option to a `my.cnf` file (thus causing the server to run as `root`), `mysqld` uses only the first `--user` option specified and produces a warning if there are multiple `--user` options. Options in `/etc/my.cnf` and `$MYSQL_HOME/my.cnf` are processed before command-line options, so it is recommended that you put a `--user` option in `/etc/my.cnf` and specify a value other than `root`. The option in `/etc/my.cnf` is found before any other `--user` options, which ensures that the server runs as a user other than `root`, and that a warning results if any other `--user` option is found.

- `--verbose`, `-v`

  Use this option with the `--help` option for detailed help.

- `--version`, `-V`

  Display version information and exit.

5.1.5 Server System Variables

The MySQL server maintains many system variables that indicate how it is configured. Each system variable has a default value. System variables can be set at server startup using options on the command line or in an option file. Most of them can be changed dynamically at runtime using the `SET` statement, which enables you to modify operation of the server without having to stop and restart it. Setting the global value of a system variable requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege. For some system variables, setting the session value also requires the same privilege; if so, it is indicated in the variable description. You can also use system variable values in expressions.

There are several ways to see the names and values of system variables:

- To see the values that a server will use based on its compiled-in defaults and any option files that it reads, use this command:

```
mysqld --verbose --help
```

- To see the values that a server will use based on its compiled-in defaults, ignoring the settings in any option files, use this command:

```
mysqld --no-defaults --verbose --help
```

- To see the current values used by a running server, use the `SHOW VARIABLES` statement or the Performance Schema system variable tables. See Section 25.11.13, “Performance Schema System Variable Tables”.

This section includes a table that lists all system variables and following the table provides a description of each one. Variables with no version indicated are present in all MySQL 8.0 releases. For more information about manipulation of system variables, see Section 5.1.6, “Using System Variables”.

Table 5.2 System Variable Summary

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
auto_generate_certs	Yes	Yes	Yes	Global	No
auto_increment_increment			Yes	Both	Yes
auto_increment_offset			Yes	Both	Yes
autocommit	Yes	Yes	Yes	Both	Yes
automatic_sp_privileges			Yes	Global	Yes
avoid_temporal_upgrade	Yes	Yes	Yes	Global	Yes
back_log			Yes	Global	No
basedir	Yes	Yes	Yes	Global	No
big-tables	Yes	Yes			Yes
- *Variable*: big_tables			Yes	Both	Yes
bind-address	Yes	Yes			No
- *Variable*: bind_address			Yes	Global	No
binlog_cache_size	Yes	Yes	Yes	Global	Yes
binlog_checksum			Yes	Global	Yes
binlog_direct_non_transactional_updates	Yes	Yes	Yes	Both	Yes
binlog_error_action	Yes	Yes	Yes	Both	Yes
binlog_expire_logs_seconds	Yes	Yes	Yes	Global	Yes
binlog-format	Yes	Yes			Yes
- *Variable*: binlog_format			Yes	Both	Yes
binlog_group_commit_sync_delay	Yes	Yes	Yes	Global	Yes
binlog_group_commit_sync_no_delay_count	Yes	Yes	Yes	Global	Yes
binlog_gtid_simple_recovery	Yes	Yes	Yes	Global	No
binlog_max_flush_queue_time			Yes	Global	Yes
binlog_order_commits			Yes	Global	Yes
binlog_row_image	Yes	Yes	Yes	Both	Yes
binlog_row_metadata	Yes	Yes	Yes	Global	Yes
binlog_rows_query_log_events			Yes	Both	Yes
binlog_stmt_cache_size	Yes	Yes	Yes	Global	Yes
binlog_transaction_dependency_history_size	Yes	Yes	Yes	Global	Yes
binlog_transaction_dependency_tracking	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
binlogging_impossible_mode	Yes	Yes	Yes	Both	Yes
block_encryption_mode	Yes	Yes	Yes	Both	Yes
bulk_insert_buffer_size	Yes	Yes	Yes	Both	Yes
character_set_client			Yes	Both	Yes
character_set_connection			Yes	Both	Yes
character_set_database[a]			Yes	Both	Yes
character-set-filesystem	Yes	Yes			Yes
- *Variable*: character_set_filesystem			Yes	Both	Yes
character_set_results			Yes	Both	Yes
character-set-server	Yes	Yes			Yes
- *Variable*: character_set_server			Yes	Both	Yes
character_set_system			Yes	Global	No
character-sets-dir	Yes	Yes			No
- *Variable*: character_sets_dir			Yes	Global	No
check_proxy_users	Yes	Yes	Yes	Global	Yes
collation_connection			Yes	Both	Yes
collation_database[b]			Yes	Both	Yes
collation-server	Yes	Yes			Yes
- *Variable*: collation_server			Yes	Both	Yes
completion_type	Yes	Yes	Yes	Both	Yes
concurrent_insert	Yes	Yes	Yes	Global	Yes
connect_timeout	Yes	Yes	Yes	Global	Yes
connection_control_failed_connections_threshold	Yes	Yes	Yes	Global	Yes
connection_control_max_connection_delay	Yes	Yes	Yes	Global	Yes
connection_control_min_connection_delay	Yes	Yes	Yes	Global	Yes
core_file			Yes	Global	No
daemon_memcached_enable_binlog	Yes	Yes	Yes	Global	No
daemon_memcached_engine_lib_name	Yes	Yes	Yes	Global	No
daemon_memcached_engine_lib_path	Yes	Yes	Yes	Global	No
daemon_memcached_option	Yes	Yes	Yes	Global	No
daemon_memcached_r_batch_size	Yes	Yes	Yes	Global	No
daemon_memcached_w_batch_size	Yes	Yes	Yes	Global	No
datadir	Yes	Yes	Yes	Global	No
date_format			Yes	Global	No
datetime_format			Yes	Global	No
debug	Yes	Yes	Yes	Both	Yes
debug_sync			Yes	Session	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
default_authentication_plugin	Yes	Yes	Yes	Global	No
default_password_lifetime	Yes	Yes	Yes	Global	Yes
default-storage-engine	Yes	Yes			Yes
- *Variable*: default_storage_engine			Yes	Both	Yes
default_tmp_storage_engine	Yes	Yes	Yes	Both	Yes
default_week_format	Yes	Yes	Yes	Both	Yes
delay-key-write	Yes	Yes			Yes
- *Variable*: delay_key_write			Yes	Global	Yes
delayed_insert_limit	Yes	Yes	Yes	Global	Yes
delayed_insert_timeout	Yes	Yes	Yes	Global	Yes
delayed_queue_size	Yes	Yes	Yes	Global	Yes
disabled_storage_engines	Yes	Yes	Yes	Global	No
disconnect_on_expired_password	Yes	Yes	Yes	Session	No
div_precision_increment	Yes	Yes	Yes	Both	Yes
end_markers_in_json			Yes	Both	Yes
enforce-gtid-consistency	Yes	Yes	Yes	Global	Yes
enforce_gtid_consistency	Yes	Yes	Yes	Global	Yes
eq_range_index_dive_limit			Yes	Both	Yes
error_count			Yes	Session	No
event-scheduler	Yes	Yes			Yes
- *Variable*: event_scheduler			Yes	Global	Yes
executed_gtids_compression_period			Yes	Global	Yes
expire_logs_days	Yes	Yes	Yes	Global	Yes
explicit_defaults_for_timestamp	Yes	Yes	Yes	Both	Yes
external_user			Yes	Session	No
flush	Yes	Yes	Yes	Global	Yes
flush_time	Yes	Yes	Yes	Global	Yes
foreign_key_checks			Yes	Both	Yes
ft_boolean_syntax	Yes	Yes	Yes	Global	Yes
ft_max_word_len	Yes	Yes	Yes	Global	No
ft_min_word_len	Yes	Yes	Yes	Global	No
ft_query_expansion_limit	Yes	Yes	Yes	Global	No
ft_stopword_file	Yes	Yes	Yes	Global	No
general-log	Yes	Yes			Yes
- *Variable*: general_log			Yes	Global	Yes
general_log_file	Yes	Yes	Yes	Global	Yes
group_concat_max_len	Yes	Yes	Yes	Both	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
group_replication_allow_local_disjoint_gtids_join	Yes	Yes	Yes	Global	Yes
group_replication_allow_local_lower_version_join	Yes	Yes	Yes	Global	Yes
group_replication_auto_increment_increment	Yes	Yes	Yes	Global	Yes
group_replication_bootstrap_group	Yes	Yes	Yes	Global	Yes
group_replication_components_stop_timeout	Yes	Yes	Yes	Global	Yes
group_replication_compression_threshold	Yes	Yes	Yes	Global	Yes
group_replication_enforce_update_everywhere_checks	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_applier_threshold	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_certifier_threshold	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_hold_percent	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_max_commit_quota	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_member_quota_percent	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_min_quota	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_min_recovery_quota	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_mode	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_period	Yes	Yes	Yes	Global	Yes
group_replication_flow_control_release_percent	Yes	Yes	Yes	Global	Yes
group_replication_force_members	Yes	Yes	Yes	Global	Yes
group_replication_group_name	Yes	Yes	Yes	Global	Yes
group_replication_group_seeds	Yes	Yes	Yes	Global	Yes
group_replication_gtid_assignment_block_size	Yes	Yes	Yes	Global	Yes
group_replication_ip_whitelist	Yes	Yes	Yes	Global	Yes
group_replication_local_address	Yes	Yes	Yes	Global	Yes
group_replication_member_weight	Yes	Yes	Yes	Global	Yes
group_replication_poll_spin_loops	Yes	Yes	Yes	Global	Yes
group_replication_recovery_complete_at	Yes	Yes	Yes	Global	Yes
group_replication_recovery_reconnect_interval	Yes	Yes	Yes	Global	Yes
group_replication_recovery_retry_count	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_ca	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_capath	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_cert	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_cipher	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_crl	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_crlpath	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_key	Yes	Yes	Yes	Global	Yes
group_replication_recovery_ssl_verify_server_cert	Yes	Yes	Yes	Global	Yes
group_replication_recovery_use_ssl	Yes	Yes	Yes	Global	Yes
group_replication_single_primary_mode	Yes	Yes	Yes	Global	Yes
group_replication_ssl_mode	Yes	Yes	Yes	Global	Yes
group_replication_start_on_boot	Yes	Yes	Yes	Global	Yes
group_replication_transaction_size_limit	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
group_replication_unreachable_majority_timeout	Yes	Yes	Yes	Global	Yes
gtid_executed			Yes	Varies	No
gtid_executed_compression_period			Yes	Global	Yes
gtid-mode	Yes	Yes			Yes
- *Variable*: gtid_mode			Yes	Global	Yes
gtid_mode			Yes	Global	Yes
gtid_next			Yes	Session	Yes
gtid_owned			Yes	Both	No
gtid_purged			Yes	Global	Yes
have_compress			Yes	Global	No
have_crypt			Yes	Global	No
have_dynamic_loading			Yes	Global	No
have_geometry			Yes	Global	No
have_openssl			Yes	Global	No
have_profiling			Yes	Global	No
have_query_cache			Yes	Global	No
have_rtree_keys			Yes	Global	No
have_ssl			Yes	Global	No
have_statement_timeout			Yes	Global	No
have_symlink			Yes	Global	No
histogram_generation_max_mem_size	Yes	Yes	Yes	Both	Yes
host_cache_size			Yes	Global	Yes
hostname			Yes	Global	No
identity			Yes	Session	Yes
ignore-builtin-innodb	Yes	Yes			No
- *Variable*: ignore_builtin_innodb			Yes	Global	No
information_schema_stats	Yes	Yes	Yes	Both	Yes
init_connect	Yes	Yes	Yes	Global	Yes
init-file	Yes	Yes			No
- *Variable*: init_file			Yes	Global	No
init_slave	Yes	Yes	Yes	Global	Yes
innodb_adaptive_flushing	Yes	Yes	Yes	Global	Yes
innodb_adaptive_flushing_lwm	Yes	Yes	Yes	Global	Yes
innodb_adaptive_hash_index	Yes	Yes	Yes	Global	Yes
innodb_adaptive_hash_index_parts	Yes	Yes	Yes	Global	No
innodb_adaptive_max_sleep_delay	Yes	Yes	Yes	Global	Yes
innodb_api_bk_commit_interval	Yes	Yes	Yes	Global	Yes
innodb_api_disable_rowlock	Yes	Yes	Yes	Global	No
innodb_api_enable_binlog	Yes	Yes	Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
innodb_api_enable_mdl	Yes	Yes	Yes	Global	No
innodb_api_trx_level	Yes	Yes	Yes	Global	Yes
innodb_autoextend_increment	Yes	Yes	Yes	Global	Yes
innodb_autoinc_lock_mode	Yes	Yes	Yes	Global	No
innodb_background_drop_list_empty	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_chunk_size	Yes	Yes	Yes	Global	No
innodb_buffer_pool_debug	Yes	Yes	Yes	Global	No
innodb_buffer_pool_dump_at_shutdown	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_dump_now	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_dump_pct	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_filename	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_instances	Yes	Yes	Yes	Global	No
innodb_buffer_pool_load_abort	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_load_at_startup	Yes	Yes	Yes	Global	No
innodb_buffer_pool_load_now	Yes	Yes	Yes	Global	Yes
innodb_buffer_pool_size	Yes	Yes	Yes	Global	Yes
innodb_change_buffer_max_size	Yes	Yes	Yes	Global	Yes
innodb_change_buffering	Yes	Yes	Yes	Global	Yes
innodb_change_buffering_debug	Yes	Yes	Yes	Global	Yes
innodb_checkpoint_disabled	Yes	Yes	Yes	Global	Yes
innodb_checksum_algorithm	Yes	Yes	Yes	Global	Yes
innodb_cmp_per_index_enabled	Yes	Yes	Yes	Global	Yes
innodb_commit_concurrency	Yes	Yes	Yes	Global	Yes
innodb_compress_debug	Yes	Yes	Yes	Global	Yes
innodb_compression_failure_threshold_pct	Yes	Yes	Yes	Global	Yes
innodb_compression_level	Yes	Yes	Yes	Global	Yes
innodb_compression_pad_pct_max	Yes	Yes	Yes	Global	Yes
innodb_concurrency_tickets	Yes	Yes	Yes	Global	Yes
innodb_data_file_path	Yes	Yes	Yes	Global	No
innodb_data_home_dir	Yes	Yes	Yes	Global	No
innodb_deadlock_detect	Yes	Yes	Yes	Global	Yes
innodb_default_row_format	Yes	Yes	Yes	Global	Yes
innodb_disable_sort_file_cache	Yes	Yes	Yes	Global	Yes
innodb_doublewrite	Yes	Yes	Yes	Global	No
innodb_fast_shutdown	Yes	Yes	Yes	Global	Yes
innodb_fil_make_page_dirty_debug	Yes	Yes	Yes	Global	Yes
innodb_file_per_table	Yes	Yes	Yes	Global	Yes
innodb_fill_factor	Yes	Yes	Yes	Global	Yes
innodb_flush_log_at_timeout			Yes	Global	Yes
innodb_flush_log_at_trx_commit	Yes	Yes	Yes	Global	Yes
innodb_flush_method	Yes	Yes	Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
innodb_flush_neighbors	Yes	Yes	Yes	Global	Yes
innodb_flush_sync	Yes	Yes	Yes	Global	Yes
innodb_flushing_avg_loops	Yes	Yes	Yes	Global	Yes
innodb_force_load_corrupted	Yes	Yes	Yes	Global	No
innodb_force_recovery	Yes	Yes	Yes	Global	No
innodb_ft_aux_table	Yes	Yes	Yes	Global	Yes
innodb_ft_cache_size	Yes	Yes	Yes	Global	No
innodb_ft_enable_diag_print	Yes	Yes	Yes	Global	Yes
innodb_ft_enable_stopword	Yes	Yes	Yes	Global	Yes
innodb_ft_max_token_size	Yes	Yes	Yes	Global	No
innodb_ft_min_token_size	Yes	Yes	Yes	Global	No
innodb_ft_num_word_optimize	Yes	Yes	Yes	Global	Yes
innodb_ft_result_cache_limit	Yes	Yes	Yes	Global	Yes
innodb_ft_server_stopword_table	Yes	Yes	Yes	Global	Yes
innodb_ft_sort_pll_degree	Yes	Yes	Yes	Global	No
innodb_ft_total_cache_size	Yes	Yes	Yes	Global	No
innodb_ft_user_stopword_table	Yes	Yes	Yes	Both	Yes
innodb_io_capacity	Yes	Yes	Yes	Global	Yes
innodb_io_capacity_max	Yes	Yes	Yes	Global	Yes
innodb_limit_optimistic_insert_debug	Yes	Yes	Yes	Global	Yes
innodb_lock_wait_timeout	Yes	Yes	Yes	Both	Yes
innodb_log_buffer_size	Yes	Yes	Yes	Global	No
innodb_log_checksums	Yes	Yes	Yes	Global	Yes
innodb_log_compressed_pages	Yes	Yes	Yes	Global	Yes
innodb_log_file_size	Yes	Yes	Yes	Global	No
innodb_log_files_in_group	Yes	Yes	Yes	Global	No
innodb_log_group_home_dir	Yes	Yes	Yes	Global	No
innodb_log_write_ahead_size	Yes	Yes	Yes	Global	Yes
innodb_lru_scan_depth	Yes	Yes	Yes	Global	Yes
innodb_max_dirty_pages_pct	Yes	Yes	Yes	Global	Yes
innodb_max_dirty_pages_pct_lwm	Yes	Yes	Yes	Global	Yes
innodb_max_purge_lag	Yes	Yes	Yes	Global	Yes
innodb_max_purge_lag_delay	Yes	Yes	Yes	Global	Yes
innodb_max_undo_log_size	Yes	Yes	Yes	Global	Yes
innodb_merge_threshold_set_all_debug	Yes	Yes	Yes	Global	Yes
innodb_monitor_disable	Yes	Yes	Yes	Global	Yes
innodb_monitor_enable	Yes	Yes	Yes	Global	Yes
innodb_monitor_reset	Yes	Yes	Yes	Global	Yes
innodb_monitor_reset_all	Yes	Yes	Yes	Global	Yes
innodb_numa_interleave	Yes	Yes	Yes	Global	No
innodb_old_blocks_pct	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
innodb_old_blocks_time	Yes	Yes	Yes	Global	Yes
innodb_online_alter_log_max_size	Yes	Yes	Yes	Global	Yes
innodb_open_files	Yes	Yes	Yes	Global	No
innodb_optimize_fulltext_only	Yes	Yes	Yes	Global	Yes
innodb_page_cleaners	Yes	Yes	Yes	Global	No
innodb_page_size	Yes	Yes	Yes	Global	No
innodb_print_all_deadlocks	Yes	Yes	Yes	Global	Yes
innodb_purge_batch_size	Yes	Yes	Yes	Global	Yes
innodb_purge_rseg_truncate_frequency	Yes	Yes	Yes	Global	Yes
innodb_purge_threads	Yes	Yes	Yes	Global	No
innodb_random_read_ahead	Yes	Yes	Yes	Global	Yes
innodb_read_ahead_threshold	Yes	Yes	Yes	Global	Yes
innodb_read_io_threads	Yes	Yes	Yes	Global	No
innodb_read_only	Yes	Yes	Yes	Global	No
innodb_redo_log_encrypt	Yes	Yes	Yes	Global	Yes
innodb_replication_delay	Yes	Yes	Yes	Global	Yes
innodb_rollback_on_timeout	Yes	Yes	Yes	Global	No
innodb_rollback_segments	Yes	Yes	Yes	Global	Yes
innodb_saved_page_number_debug	Yes	Yes	Yes	Global	Yes
innodb_scan_directories	Yes	Yes	Yes	Global	No
innodb_sort_buffer_size	Yes	Yes	Yes	Global	No
innodb_spin_wait_delay	Yes	Yes	Yes	Global	Yes
innodb_stats_auto_recalc	Yes	Yes	Yes	Global	Yes
innodb_stats_include_delete_marked	Yes	Yes	Yes	Global	Yes
innodb_stats_method	Yes	Yes	Yes	Global	Yes
innodb_stats_on_metadata	Yes	Yes	Yes	Global	Yes
innodb_stats_persistent	Yes	Yes	Yes	Global	Yes
innodb_stats_persistent_sample_pages	Yes	Yes	Yes	Global	Yes
innodb_stats_transient_sample_pages	Yes	Yes	Yes	Global	Yes
innodb_status_output	Yes	Yes	Yes	Global	Yes
innodb_status_output_locks	Yes	Yes	Yes	Global	Yes
innodb_strict_mode	Yes	Yes	Yes	Both	Yes
innodb_sync_array_size	Yes	Yes	Yes	Global	No
innodb_sync_debug	Yes	Yes	Yes	Global	No
innodb_sync_spin_loops	Yes	Yes	Yes	Global	Yes
innodb_table_locks	Yes	Yes	Yes	Both	Yes
innodb_temp_data_file_path	Yes	Yes	Yes	Global	No
innodb_thread_concurrency	Yes	Yes	Yes	Global	Yes
innodb_thread_sleep_delay	Yes	Yes	Yes	Global	Yes
innodb_tmpdir	Yes	Yes	Yes	Both	Yes
innodb_trx_purge_view_update_only_debug	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
innodb_trx_rseg_n_slots_debug	Yes	Yes	Yes	Global	Yes
innodb_undo_directory	Yes	Yes	Yes	Global	No
innodb_undo_log_encrypt	Yes	Yes	Yes	Global	Yes
innodb_undo_log_truncate	Yes	Yes	Yes	Global	Yes
innodb_undo_logs	Yes	Yes	Yes	Global	Yes
innodb_undo_tablespaces	Yes	Yes	Yes	Global	Varies
innodb_use_native_aio	Yes	Yes	Yes	Global	No
innodb_version			Yes	Global	No
innodb_write_io_threads	Yes	Yes	Yes	Global	No
insert_id			Yes	Session	Yes
interactive_timeout	Yes	Yes	Yes	Both	Yes
internal_tmp_disk_storage_engine	Yes	Yes	Yes	Global	Yes
internal_tmp_mem_storage_engine	Yes	Yes	Yes	Session	Yes
join_buffer_size	Yes	Yes	Yes	Both	Yes
keep_files_on_create	Yes	Yes	Yes	Both	Yes
key_buffer_size	Yes	Yes	Yes	Global	Yes
key_cache_age_threshold	Yes	Yes	Yes	Global	Yes
key_cache_block_size	Yes	Yes	Yes	Global	Yes
key_cache_division_limit	Yes	Yes	Yes	Global	Yes
keyring_file_data	Yes	Yes	Yes	Global	Yes
language	Yes	Yes	Yes	Global	No
large_files_support			Yes	Global	No
large_page_size			Yes	Global	No
large-pages	Yes	Yes			No
- *Variable*: large_pages			Yes	Global	No
last_insert_id			Yes	Session	Yes
lc-messages	Yes	Yes			Yes
- *Variable*: lc_messages			Yes	Both	Yes
lc-messages-dir	Yes	Yes			No
- *Variable*: lc_messages_dir			Yes	Global	No
lc_time_names			Yes	Both	Yes
license			Yes	Global	No
local_infile			Yes	Global	Yes
lock_wait_timeout	Yes	Yes	Yes	Both	Yes
locked_in_memory			Yes	Global	No
log-bin	Yes	Yes	Yes	Global	No
log_bin			Yes	Global	No
log_bin_basename			Yes	Global	No
log_bin_index			Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
log-bin-trust-function-creators	Yes	Yes			Yes
- *Variable*: log_bin_trust_function_creators			Yes	Global	Yes
log-bin-use-v1-row-events	Yes	Yes			No
- *Variable*: log_bin_use_v1_row_events			Yes	Global	No
log_bin_use_v1_row_events	Yes	Yes	Yes	Global	No
log_builtin_as_identified_by_password	Yes	Yes	Yes	Global	Yes
log-error	Yes	Yes			No
- *Variable*: log_error			Yes	Global	No
log_error_verbosity	Yes	Yes	Yes	Global	Yes
log-output	Yes	Yes			Yes
- *Variable*: log_output			Yes	Global	Yes
log-queries-not-using-indexes	Yes	Yes			Yes
- *Variable*: log_queries_not_using_indexes			Yes	Global	Yes
log-slave-updates	Yes	Yes			No
- *Variable*: log_slave_updates			Yes	Global	No
log_slave_updates	Yes	Yes	Yes	Global	No
log_slow_admin_statements			Yes	Global	Yes
log_slow_slave_statements			Yes	Global	Yes
log_statements_unsafe_for_binlog			Yes	Global	Yes
log_syslog	Yes	Yes	Yes	Global	Yes
log_syslog_facility	Yes	Yes	Yes	Global	Yes
log_syslog_include_pid	Yes	Yes	Yes	Global	Yes
log_syslog_tag	Yes	Yes	Yes	Global	Yes
log_throttle_queries_not_using_indexes			Yes	Global	Yes
log_timestamps	Yes	Yes	Yes	Global	Yes
log-warnings	Yes	Yes			Yes
- *Variable*: log_warnings			Yes	Global	Yes
long_query_time	Yes	Yes	Yes	Both	Yes
low-priority-updates	Yes	Yes			Yes
- *Variable*: low_priority_updates			Yes	Both	Yes
lower_case_file_system			Yes	Global	No
lower_case_table_names	Yes	Yes	Yes	Global	No
master_info_repository	Yes	Yes	Yes	Global	Yes
master_verify_checksum			Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
max_allowed_packet	Yes	Yes	Yes	Both	Yes
max_binlog_cache_size	Yes	Yes	Yes	Global	Yes
max_binlog_size	Yes	Yes	Yes	Global	Yes
max_binlog_stmt_cache_size	Yes	Yes	Yes	Global	Yes
max_connect_errors	Yes	Yes	Yes	Global	Yes
max_connections	Yes	Yes	Yes	Global	Yes
max_delayed_threads	Yes	Yes	Yes	Both	Yes
max_digest_length	Yes	Yes	Yes	Global	No
max_error_count	Yes	Yes	Yes	Both	Yes
max_execution_time	Yes	Yes	Yes	Both	Yes
max_heap_table_size	Yes	Yes	Yes	Both	Yes
max_insert_delayed_threads			Yes	Both	Yes
max_join_size	Yes	Yes	Yes	Both	Yes
max_length_for_sort_data	Yes	Yes	Yes	Both	Yes
max_points_in_geometry	Yes	Yes	Yes	Global	Yes
max_prepared_stmt_count	Yes	Yes	Yes	Global	Yes
max_relay_log_size	Yes	Yes	Yes	Global	Yes
max_seeks_for_key	Yes	Yes	Yes	Both	Yes
max_sort_length	Yes	Yes	Yes	Both	Yes
max_sp_recursion_depth	Yes	Yes	Yes	Both	Yes
max_tmp_tables			Yes	Both	Yes
max_user_connections	Yes	Yes	Yes	Both	Yes
max_write_lock_count	Yes	Yes	Yes	Global	Yes
mecab_rc_file	Yes	Yes	Yes	Global	No
metadata_locks_cache_size			Yes	Global	No
metadata_locks_hash_instances			Yes	Global	No
min-examined-row-limit	Yes	Yes	Yes	Both	Yes
myisam_data_pointer_size	Yes	Yes	Yes	Global	Yes
myisam_max_sort_file_size	Yes	Yes	Yes	Global	Yes
myisam_mmap_size	Yes	Yes	Yes	Global	No
myisam_recover_options			Yes	Global	No
myisam_repair_threads	Yes	Yes	Yes	Both	Yes
myisam_sort_buffer_size	Yes	Yes	Yes	Both	Yes
myisam_stats_method	Yes	Yes	Yes	Both	Yes
myisam_use_mmap	Yes	Yes	Yes	Global	Yes
mysql_native_password_proxy_users	Yes	Yes	Yes	Global	Yes
mysqlx_bind_address	Yes	Yes	Yes	Global	No
mysqlx_connect_timeout	Yes	Yes	Yes	Global	Yes
mysqlx_max_connections	Yes	Yes	Yes	Global	Yes
mysqlx_port	Yes	Yes	Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
mysqlx_port_open_timeout	Yes	Yes	Yes	Global	No
mysqlx-socket	Yes	Yes	Yes	Global	No
mysqlx_socket	Yes	Yes	Yes	Global	No
mysqlx_ssl	Yes	Yes	Yes	Global	No
mysqlx_ssl_ca	Yes	Yes	Yes	Global	No
mysqlx_ssl_capath	Yes	Yes	Yes	Global	No
mysqlx_ssl_cert	Yes	Yes	Yes	Global	No
mysqlx_ssl_crl	Yes	Yes	Yes	Global	No
mysqlx_ssl_crlpath	Yes	Yes	Yes	Global	No
mysqlx_ssl_key	Yes	Yes	Yes	Global	No
named_pipe			Yes	Global	No
ndb-batch-size	Yes	Yes	Yes	Global	No
ndb-blob-write-batch-bytes	Yes	Yes	Yes	Both	Yes
ndb-cluster-connection-pool	Yes	Yes	Yes	Global	No
ndb-cluster-connection-pool-nodeids	Yes	Yes	Yes	Global	No
ndb-deferred-constraints	Yes	Yes			Yes
- *Variable*: ndb_deferred_constraints			Yes	Both	Yes
ndb_deferred_constraints	Yes	Yes	Yes	Both	Yes
ndb-distribution	Yes	Yes			Yes
- *Variable*: ndb_distribution			Yes	Global	Yes
ndb_distribution	Yes	Yes	Yes	Global	Yes
ndb_eventbuffer_free_percent	Yes	Yes	Yes	Global	Yes
ndb_eventbuffer_max_alloc	Yes	Yes	Yes	Global	Yes
ndb_force_send	Yes	Yes	Yes	Both	Yes
ndb_index_stat_enable	Yes	Yes	Yes	Both	Yes
ndb_index_stat_option	Yes	Yes	Yes	Both	Yes
ndb_join_pushdown			Yes	Both	Yes
ndb-log-apply-status	Yes	Yes			No
- *Variable*: ndb_log_apply_status			Yes	Global	No
ndb_log_apply_status	Yes	Yes	Yes	Global	No
ndb_log_binlog_index	Yes		Yes	Global	Yes
ndb-log-empty-epochs	Yes	Yes	Yes	Global	Yes
ndb-log-empty-update	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
ndb-log-transaction-id	Yes	Yes			No
- *Variable*: ndb_log_transaction_id			Yes	Global	No
ndb_log_updated_only	Yes	Yes	Yes	Global	Yes
ndb_optimization_delay			Yes	Global	Yes
ndb_optimized_node_selection	Yes	Yes	Yes	Global	No
ndb_recv_thread_cpu_mask			Yes	Global	Yes
ndb_report_thresh_binlog_epoch_slip	Yes	Yes	Yes	Global	Yes
ndb_report_thresh_binlog_mem_usage	Yes	Yes	Yes	Global	Yes
ndb_show_foreign_key_mock_tables	Yes	Yes	Yes	Global	Yes
Ndb_slave_max_replicated_epoch			Yes	Global	No
ndb_table_no_logging			Yes	Session	Yes
ndb_use_transactions	Yes	Yes	Yes	Both	Yes
ndb_version			Yes	Global	No
ndb_version_string			Yes	Global	No
ndb-wait-setup	Yes	Yes	Yes	Global	No
ndbinfo_database			Yes	Global	No
ndbinfo_max_rows	Yes		Yes	Both	Yes
ndbinfo_show_hidden	Yes		Yes	Both	Yes
ndbinfo_version			Yes	Global	No
net_buffer_length	Yes	Yes	Yes	Both	Yes
net_read_timeout	Yes	Yes	Yes	Both	Yes
net_retry_count	Yes	Yes	Yes	Both	Yes
net_write_timeout	Yes	Yes	Yes	Both	Yes
new	Yes	Yes	Yes	Both	Yes
ngram_token_size	Yes	Yes	Yes	Global	No
offline_mode	Yes	Yes	Yes	Global	Yes
old	Yes	Yes	Yes	Global	No
old-alter-table	Yes	Yes			Yes
- *Variable*: old_alter_table			Yes	Both	Yes
old_passwords			Yes	Both	Yes
open-files-limit	Yes	Yes			No
- *Variable*: open_files_limit			Yes	Global	No
optimizer_prune_level	Yes	Yes	Yes	Both	Yes
optimizer_search_depth	Yes	Yes	Yes	Both	Yes
optimizer_switch	Yes	Yes	Yes	Both	Yes
optimizer_trace			Yes	Both	Yes
optimizer_trace_features			Yes	Both	Yes
optimizer_trace_limit			Yes	Both	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
optimizer_trace_max_mem_size			Yes	Both	Yes
optimizer_trace_offset			Yes	Both	Yes
parser_max_mem_size	Yes	Yes	Yes	Both	Yes
performance_schema	Yes	Yes	Yes	Global	No
performance_schema_accounts_size	Yes	Yes	Yes	Global	No
performance_schema_digests_size	Yes	Yes	Yes	Global	No
performance_schema_error_size	Yes	Yes	Yes	Global	No
performance_schema_events_stages_history_long_size	Yes	Yes	Yes	Global	No
performance_schema_events_stages_history_size	Yes	Yes	Yes	Global	No
performance_schema_events_statements_history_long_size	Yes	Yes	Yes	Global	No
performance_schema_events_statements_history_size	Yes	Yes	Yes	Global	No
performance_schema_events_transactions_history_long_size	Yes	Yes	Yes	Global	No
performance_schema_events_transactions_history_size	Yes	Yes	Yes	Global	No
performance_schema_events_waits_history_long_size	Yes	Yes	Yes	Global	No
performance_schema_events_waits_history_size	Yes	Yes	Yes	Global	No
performance_schema_hosts_size	Yes	Yes	Yes	Global	No
performance_schema_max_cond_classes	Yes	Yes	Yes	Global	No
performance_schema_max_cond_instances	Yes	Yes	Yes	Global	No
performance_schema_max_digest_length	Yes	Yes	Yes	Global	No
performance_schema_max_file_classes	Yes	Yes	Yes	Global	No
performance_schema_max_file_handles	Yes	Yes	Yes	Global	No
performance_schema_max_file_instances	Yes	Yes	Yes	Global	No
performance_schema_max_index_stat	Yes	Yes	Yes	Global	No
performance_schema_max_memory_classes	Yes	Yes	Yes	Global	No
performance_schema_max_metadata_locks	Yes	Yes	Yes	Global	No
performance_schema_max_mutex_classes	Yes	Yes	Yes	Global	No
performance_schema_max_mutex_instances	Yes	Yes	Yes	Global	No
performance_schema_max_prepared_statements_instances	Yes	Yes	Yes	Global	No
performance_schema_max_program_instances	Yes	Yes	Yes	Global	No
performance_schema_max_rwlock_classes	Yes	Yes	Yes	Global	No
performance_schema_max_rwlock_instances	Yes	Yes	Yes	Global	No
performance_schema_max_socket_classes	Yes	Yes	Yes	Global	No
performance_schema_max_socket_instances	Yes	Yes	Yes	Global	No
performance_schema_max_sql_text_length	Yes	Yes	Yes	Global	No
performance_schema_max_stage_classes	Yes	Yes	Yes	Global	No
performance_schema_max_statement_classes	Yes	Yes	Yes	Global	No
performance_schema_max_statement_stack	Yes	Yes	Yes	Global	No
performance_schema_max_table_handles	Yes	Yes	Yes	Global	No
performance_schema_max_table_instances	Yes	Yes	Yes	Global	No
performance_schema_max_table_lock_stat	Yes	Yes	Yes	Global	No
performance_schema_max_thread_classes	Yes	Yes	Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
performance_schema_max_thread_instances	Yes	Yes	Yes	Global	No
performance_schema_session_connect_attrs_size	Yes	Yes	Yes	Global	No
performance_schema_setup_actors_size	Yes	Yes	Yes	Global	No
performance_schema_setup_objects_size	Yes	Yes	Yes	Global	No
performance_schema_users_size	Yes	Yes	Yes	Global	No
persisted_globals_load	Yes	Yes	Yes	Global	No
pid-file	Yes	Yes			No
- *Variable*: pid_file			Yes	Global	No
plugin_dir	Yes	Yes	Yes	Global	No
port	Yes	Yes	Yes	Global	No
preload_buffer_size	Yes	Yes	Yes	Both	Yes
profiling			Yes	Both	Yes
profiling_history_size	Yes	Yes	Yes	Both	Yes
protocol_version			Yes	Global	No
proxy_user			Yes	Session	No
pseudo_slave_mode			Yes	Session	Yes
pseudo_thread_id			Yes	Session	Yes
query_alloc_block_size	Yes	Yes	Yes	Both	Yes
query_cache_limit	Yes	Yes	Yes	Global	Yes
query_cache_min_res_unit	Yes	Yes	Yes	Global	Yes
query_cache_size	Yes	Yes	Yes	Global	Yes
query_cache_type	Yes	Yes	Yes	Both	Yes
query_cache_wlock_invalidate	Yes	Yes	Yes	Both	Yes
query_prealloc_size	Yes	Yes	Yes	Both	Yes
rand_seed1			Yes	Session	Yes
rand_seed2			Yes	Session	Yes
range_alloc_block_size	Yes	Yes	Yes	Both	Yes
range_optimizer_max_mem_size	Yes	Yes	Yes	Both	Yes
rbr_exec_mode			Yes	Session	Yes
read_buffer_size	Yes	Yes	Yes	Both	Yes
read_only	Yes	Yes	Yes	Global	Yes
read_rnd_buffer_size	Yes	Yes	Yes	Both	Yes
relay-log	Yes	Yes			No
- *Variable*: relay_log			Yes	Global	No
relay_log_basename			Yes	Global	No
relay-log-index	Yes	Yes			No
- *Variable*: relay_log_index			Yes	Global	No
relay_log_index	Yes	Yes	Yes	Global	No
relay_log_info_file	Yes	Yes	Yes	Global	No
relay_log_info_repository			Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
relay_log_purge	Yes	Yes	Yes	Global	Yes
relay_log_recovery	Yes	Yes	Yes	Global	No
relay_log_space_limit	Yes	Yes	Yes	Global	No
report-host	Yes	Yes			No
- *Variable*: report_host			Yes	Global	No
report-password	Yes	Yes			No
- *Variable*: report_password			Yes	Global	No
report-port	Yes	Yes			No
- *Variable*: report_port			Yes	Global	No
report-user	Yes	Yes			No
- *Variable*: report_user			Yes	Global	No
require_secure_transport	Yes	Yes	Yes	Global	Yes
rewriter_enabled			Yes	Global	Yes
rewriter_verbose			Yes	Global	Yes
rpl_semi_sync_master_enabled			Yes	Global	Yes
rpl_semi_sync_master_timeout			Yes	Global	Yes
rpl_semi_sync_master_trace_level			Yes	Global	Yes
rpl_semi_sync_master_wait_for_slave_count			Yes	Global	Yes
rpl_semi_sync_master_wait_no_slave			Yes	Global	Yes
rpl_semi_sync_master_wait_point			Yes	Global	Yes
rpl_semi_sync_slave_enabled			Yes	Global	Yes
rpl_semi_sync_slave_trace_level			Yes	Global	Yes
rpl_stop_slave_timeout	Yes	Yes	Yes	Global	Yes
schema_definition_cache	Yes	Yes	Yes	Global	Yes
secure-auth	Yes	Yes			Yes
- *Variable*: secure_auth			Yes	Global	Yes
secure-file-priv	Yes	Yes			No
- *Variable*: secure_file_priv			Yes	Global	No
server-id	Yes	Yes			Yes
- *Variable*: server_id			Yes	Global	Yes
server_uuid			Yes	Global	No
session_track_gtids	Yes	Yes	Yes	Both	Yes
session_track_schema	Yes	Yes	Yes	Both	Yes
session_track_state_change	Yes	Yes	Yes	Both	Yes
session_track_system_variables	Yes	Yes	Yes	Both	Yes
sha256_password_auto_generate_rsa_keys	Yes	Yes	Yes	Global	No
sha256_password_private_key_path			Yes	Global	No

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
sha256_password_proxy_users	Yes	Yes	Yes	Global	Yes
sha256_password_public_key_path			Yes	Global	No
shared_memory	Yes	Yes	Yes	Global	No
shared_memory_base_name	Yes	Yes	Yes	Global	No
show_compatibility_56	Yes	Yes	Yes	Global	Yes
show_old_temporals	Yes	Yes	Yes	Both	Yes
simplified_binlog_gtid_recovery	Yes	Yes	Yes	Global	No
skip_external_locking	Yes	Yes	Yes	Global	No
skip-name-resolve	Yes	Yes			No
- *Variable*: skip_name_resolve			Yes	Global	No
skip-networking	Yes	Yes			No
- *Variable*: skip_networking			Yes	Global	No
skip-show-database	Yes	Yes			No
- *Variable*: skip_show_database			Yes	Global	No
slave_allow_batching	Yes	Yes	Yes	Global	Yes
slave_checkpoint_group	Yes	Yes	Yes	Global	Yes
slave_checkpoint_period	Yes	Yes	Yes	Global	Yes
slave_compressed_protocol	Yes	Yes	Yes	Global	Yes
slave_exec_mode	Yes	Yes	Yes	Global	Yes
slave-load-tmpdir	Yes	Yes			No
- *Variable*: slave_load_tmpdir			Yes	Global	No
slave_max_allowed_packet			Yes	Global	Yes
slave-net-timeout	Yes	Yes			Yes
- *Variable*: slave_net_timeout			Yes	Global	Yes
slave_parallel_type			Yes	Global	Yes
slave_parallel_workers	Yes		Yes	Global	Yes
slave_pending_jobs_size_max			Yes	Global	Yes
slave_preserve_commit_order	Yes		Yes	Global	Yes
slave_rows_search_algorithms			Yes	Global	Yes
slave-skip-errors	Yes	Yes			No
- *Variable*: slave_skip_errors			Yes	Global	No
slave_sql_verify_checksum			Yes	Global	Yes
slave_transaction_retries	Yes	Yes	Yes	Global	Yes
slave_type_conversions	Yes	Yes	Yes	Global	No
slow_launch_time	Yes	Yes	Yes	Global	Yes
slow-query-log	Yes	Yes			Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
- *Variable*: slow_query_log			Yes	Global	Yes
slow_query_log_file	Yes	Yes	Yes	Global	Yes
socket	Yes	Yes	Yes	Global	No
sort_buffer_size	Yes	Yes	Yes	Both	Yes
sql_auto_is_null			Yes	Both	Yes
sql_big_selects			Yes	Both	Yes
sql_buffer_result			Yes	Both	Yes
sql_log_bin			Yes	Session	Yes
sql_log_off			Yes	Both	Yes
sql-mode	Yes	Yes			Yes
- *Variable*: sql_mode			Yes	Both	Yes
sql_notes			Yes	Both	Yes
sql_quote_show_create			Yes	Both	Yes
sql_safe_updates			Yes	Both	Yes
sql_select_limit			Yes	Both	Yes
sql_slave_skip_counter			Yes	Global	Yes
sql_warnings			Yes	Both	Yes
ssl-ca	Yes	Yes			No
- *Variable*: ssl_ca			Yes	Global	No
ssl-capath	Yes	Yes			No
- *Variable*: ssl_capath			Yes	Global	No
ssl-cert	Yes	Yes			No
- *Variable*: ssl_cert			Yes	Global	No
ssl-cipher	Yes	Yes			No
- *Variable*: ssl_cipher			Yes	Global	No
ssl-crl	Yes	Yes			No
- *Variable*: ssl_crl			Yes	Global	No
ssl-crlpath	Yes	Yes			No
- *Variable*: ssl_crlpath			Yes	Global	No
ssl-key	Yes	Yes			No
- *Variable*: ssl_key			Yes	Global	No
stored_program_cache	Yes	Yes	Yes	Global	Yes
stored_program_definition_cache	Yes	Yes	Yes	Global	Yes
super_read_only	Yes	Yes	Yes	Global	Yes
sync_binlog	Yes	Yes	Yes	Global	Yes
sync_master_info	Yes	Yes	Yes	Global	Yes
sync_relay_log	Yes	Yes	Yes	Global	Yes
sync_relay_log_info	Yes	Yes	Yes	Global	Yes

Name	Cmd-Line	Option File	System Var	Var Scope	Dynamic
system_time_zone			Yes	Global	No
table_definition_cache			Yes	Global	Yes
table_open_cache			Yes	Global	Yes
table_open_cache_instances			Yes	Global	No
tablespace_definition_cache	Yes	Yes	Yes	Global	Yes
temptable_max_ram	Yes	Yes	Yes	Global	Yes
thread_cache_size	Yes	Yes	Yes	Global	Yes
thread_handling	Yes	Yes	Yes	Global	No
thread_stack	Yes	Yes	Yes	Global	No
time_format			Yes	Global	No
time_zone			Yes	Both	Yes
timestamp			Yes	Session	Yes
tls_version	Yes	Yes	Yes	Global	No
tmp_table_size	Yes	Yes	Yes	Both	Yes
tmpdir	Yes	Yes	Yes	Global	No
transaction_alloc_block_size	Yes	Yes	Yes	Both	Yes
transaction_allow_batching			Yes	Session	Yes
transaction_prealloc_size	Yes	Yes	Yes	Both	Yes
transaction_write_set_extraction	Yes		Yes	Both	Yes
tx_isolation			Yes	Both	Yes
tx_read_only			Yes	Both	Yes
unique_checks			Yes	Both	Yes
updatable_views_with_limit	Yes	Yes	Yes	Both	Yes
validate_password_check_user_name	Yes	Yes	Yes	Global	Yes
validate_password_dictionary_file			Yes	Global	Yes
validate_password_length			Yes	Global	Yes
validate_password_mixed_case_count			Yes	Global	Yes
validate_password_number_count			Yes	Global	Yes
validate_password_policy			Yes	Global	Yes
validate_password_special_char_count			Yes	Global	Yes
validate_user_plugins			Yes	Global	No
version			Yes	Global	No
version_comment			Yes	Global	No
version_compile_machine			Yes	Global	No
version_compile_os			Yes	Global	No
version_tokens_session	Yes	Yes	Yes	Both	Yes
version_tokens_session_number	Yes	Yes	Yes	Both	No
wait_timeout	Yes	Yes	Yes	Both	Yes
warning_count			Yes	Session	No

[a] This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.

[b] This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.

For additional system variable information, see these sections:

- Section 5.1.6, “Using System Variables”, discusses the syntax for setting and displaying system variable values.
- Section 5.1.6.2, “Dynamic System Variables”, lists the variables that can be set at runtime.
- Information on tuning system variables can be found in Section 5.1.1, “Configuring the Server”.
- Section 15.13, “InnoDB Startup Options and System Variables”, lists `InnoDB` system variables.
- For information on server system variables specific to replication, see Section 18.1.6, “Replication and Binary Logging Options and Variables”.

Note

Some of the following variable descriptions refer to “enabling” or “disabling” a variable. These variables can be enabled with the `SET` statement by setting them to `ON` or `1`, or disabled by setting them to `OFF` or `0`. Boolean variables can be set at startup to the values `ON`, `TRUE`, `OFF`, and `FALSE` (not case sensitive), as well as `1` and `0`. See Section 4.2.5, “Program Option Modifiers”.

Some system variables control the size of buffers or caches. For a given buffer, the server might need to allocate internal data structures. These structures typically are allocated from the total memory allocated to the buffer, and the amount of space required might be platform dependent. This means that when you assign a value to a system variable that controls a buffer size, the amount of space actually available might differ from the value assigned. In some cases, the amount might be less than the value assigned. It is also possible that the server will adjust a value upward. For example, if you assign a value of 0 to a variable for which the minimal value is 1024, the server will set the value to 1024.

Values for buffer sizes, lengths, and stack sizes are given in bytes unless otherwise specified.

Some system variables take file name values. Unless otherwise specified, the default file location is the data directory if the value is a relative path name. To specify the location explicitly, use an absolute path name. Suppose that the data directory is `/var/mysql/data`. If a file-valued variable is given as a relative path name, it will be located under `/var/mysql/data`. If the value is an absolute path name, its location is as given by the path name.

- `autocommit`

Command-Line Format	`--autocommit[=#]`	
System Variable	**Name**	`autocommit`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

The autocommit mode. If set to 1, all changes to a table take effect immediately. If set to 0, you must use `COMMIT` to accept a transaction or `ROLLBACK` to cancel it. If `autocommit` is 0 and you change it to 1, MySQL performs an automatic `COMMIT` of any open transaction. Another way to begin a transaction is to use a `START TRANSACTION` or `BEGIN` statement. See Section 13.3.1, “START TRANSACTION, COMMIT, and ROLLBACK Syntax”.

By default, client connections begin with `autocommit` set to 1. To cause clients to begin with a default of 0, set the global `autocommit` value by starting the server with the `--autocommit=0` option. To set the variable using an option file, include these lines:

```
[mysqld]
autocommit=0
```

- `automatic_sp_privileges`

System Variable	Name	automatic_sp_privileges
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	TRUE

When this variable has a value of 1 (the default), the server automatically grants the `EXECUTE` and `ALTER ROUTINE` privileges to the creator of a stored routine, if the user cannot already execute and alter or drop the routine. (The `ALTER ROUTINE` privilege is required to drop the routine.) The server also automatically drops those privileges from the creator when the routine is dropped. If `automatic_sp_privileges` is 0, the server does not automatically add or drop these privileges.

The creator of a routine is the account used to execute the `CREATE` statement for it. This might not be the same as the account named as the `DEFINER` in the routine definition.

See also Section 23.2.2, “Stored Routines and MySQL Privileges”.

- `auto_generate_certs`

Command-Line Format	--auto-generate-certs[={OFF|ON}]	
System Variable	Name	auto_generate_certs
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	ON

This variable is available if the server was compiled using OpenSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). It controls whether the server autogenerates SSL key and certificate files in the data directory, if they do not already exist.

At startup, the server automatically generates server-side and client-side SSL certificate and key files in the data directory if the `auto_generate_certs` system variable is enabled, no SSL options other than `--ssl` are specified, and the server-side SSL files are missing from the data directory. These files enable secure client connections using SSL; see Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

For more information about SSL file autogeneration, including file names and characteristics, see Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”

The `sha256_password_auto_generate_rsa_keys` system variable is related but controls autogeneration of RSA key-pair files needed for secure password exchange using RSA over unencypted connections.

- `avoid_temporal_upgrade`

Deprecated	5.7.6

Command-Line Format	--avoid-temporal-upgrade={OFF|ON}	
System Variable	Name	avoid_temporal_upgrade
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

This variable controls whether ALTER TABLE implicitly upgrades temporal columns found to be in pre-5.6.4 format (TIME, DATETIME, and TIMESTAMP columns without support for fractional seconds precision). Upgrading such columns requires a table rebuild, which prevents any use of fast alterations that might otherwise apply to the operation to be performed.

This variable is disabled by default. Enabling it causes ALTER TABLE not to rebuild temporal columns and thereby be able to take advantage of possible fast alterations.

This variable is deprecated and will be removed in a future MySQL release.

- back_log

System Variable	Name	back_log
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	-1 (autosized)
	Min Value	1
	Max Value	65535

The number of outstanding connection requests MySQL can have. This comes into play when the main MySQL thread gets very many connection requests in a very short time. It then takes some time (although very little) for the main thread to check the connection and start a new thread. The back_log value indicates how many requests can be stacked during this short time before MySQL momentarily stops answering new requests. You need to increase this only if you expect a large number of connections in a short period of time.

In other words, this value is the size of the listen queue for incoming TCP/IP connections. Your operating system has its own limit on the size of this queue. The manual page for the Unix listen() system call should have more details. Check your OS documentation for the maximum value for this variable. back_log cannot be set higher than your operating system limit.

The default value is based on the following formula, capped to a limit of 900:

```
50 + (max_connections / 5)
```

- basedir

Command-Line Format	--basedir=dir_name	
System Variable	Name	basedir

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (<= 8.0.1)	**Type**	directory name
Permitted Values (>= 8.0.2)	**Type**	directory name
	Default	`parent of mysqld installation directory`

The path to the MySQL installation base directory.

- `big_tables`

Command-Line Format	`--big-tables`	
System Variable	**Name**	`big_tables`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

If set to 1, all temporary tables are stored on disk rather than in memory. This is a little slower, but the error `The table tbl_name is full` does not occur for `SELECT` operations that require a large temporary table. The default value for a new connection is 0 (use in-memory temporary tables). Normally, you should never need to set this variable, because in-memory tables are automatically converted to disk-based tables as required.

- `binlog_expire_logs_seconds`

Introduced	8.0.1	
Command-Line Format	`--binlog-expire-logs-seconds=#`	
System Variable	**Name**	`binlog_expire_logs_seconds`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

This variable sets the binary log expiration period in seconds.

The effects of this variable and `expire_logs_days` are cumulative; for example, if `expire_logs_days` is 1 and `binlog_expire_logs_seconds` is 43200, then the binary log is purged every 1.5 days. This produces the same result as setting `binlog_expire_logs_seconds` to 129600 and leaving `expire_logs_days` set to 0.

In MySQL 8.0, both `binlog_expire_logs_seconds` and `expire_logs_days` must be set to 0 to disable automatic purging of the binary log.

To remove binary log files manually, use the `PURGE BINARY LOGS` statement. See Section 13.4.1.1, "PURGE BINARY LOGS Syntax".

- `bind_address`

Command-Line Format	`--bind-address=addr`	
System Variable	Name	`bind_address`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	`*`

The value of the `--bind-address` option.

- `block_encryption_mode`

Command-Line Format	`--block-encryption-mode=#`	
System Variable	Name	`block_encryption_mode`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string
	Default	`aes-128-ecb`

This variable controls the block encryption mode for block-based algorithms such as AES. It affects encryption for `AES_ENCRYPT()` and `AES_DECRYPT()`.

`block_encryption_mode` takes a value in `aes-`*`keylen`*`-`*`mode`* format, where *`keylen`* is the key length in bits and *`mode`* is the encryption mode. The value is not case sensitive. Permitted *`keylen`* values are 128, 192, and 256. Permitted encryption modes depend on whether MySQL was compiled using OpenSSL or yaSSL:

- For OpenSSL, permitted *`mode`* values are: `ECB`, `CBC`, `CFB1`, `CFB8`, `CFB128`, `OFB`
- For yaSSL, permitted *`mode`* values are: `ECB`, `CBC`

For example, this statement causes the AES encryption functions to use a key length of 256 bits and the CBC mode:

```
SET block_encryption_mode = 'aes-256-cbc';
```

An error occurs for attempts to set `block_encryption_mode` to a value containing an unsupported key length or a mode that the SSL library does not support.

- `bulk_insert_buffer_size`

Command-Line Format	`--bulk-insert-buffer-size=#`	
System Variable	Name	`bulk_insert_buffer_size`

	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`8388608`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`8388608`
	Min Value	`0`
	Max Value	`18446744073709551615`

`MyISAM` uses a special tree-like cache to make bulk inserts faster for `INSERT ... SELECT`, `INSERT ... VALUES (...), (...), ...`, and `LOAD DATA INFILE` when adding data to nonempty tables. This variable limits the size of the cache tree in bytes per thread. Setting it to 0 disables this optimization. The default value is 8MB.

- `character_set_client`

System Variable	**Name**	`character_set_client`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The character set for statements that arrive from the client. The session value of this variable is set using the character set requested by the client when the client connects to the server. (Many clients support a `--default-character-set` option to enable this character set to be specified explicitly. See also Section 10.1.4, “Connection Character Sets and Collations”.) The global value of the variable is used to set the session value in cases when the client-requested value is unknown or not available, or the server is configured to ignore client requests:

- The client is from a version of MySQL older than MySQL 4.1, and thus does not request a character set.
- The client requests a character set not known to the server. For example, a Japanese-enabled client requests `sjis` when connecting to a server not configured with `sjis` support.
- `mysqld` was started with the `--skip-character-set-client-handshake` option, which causes it to ignore client character set configuration. This reproduces MySQL 4.0 behavior and is useful should you wish to upgrade the server without upgrading all the clients.

`ucs2`, `utf16`, `utf16le`, and `utf32` cannot be used as a client character set, which means that they also do not work for `SET NAMES` or `SET CHARACTER SET`.

- `character_set_connection`

System Variable	**Name**	`character_set_connection`

	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The character set used for literals that do not have a character set introducer and for number-to-string conversion. For information about introducers, see Section 10.1.3.8, “Character Set Introducers”.

- `character_set_database`

System Variable	**Name**	`character_set_database`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	**Type**	string
	Default	`latin1`
Permitted Values (>= 8.0.1)	**Type**	string
	Default	`utf8mb4`
Footnote	This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.	

The character set used by the default database. The server sets this variable whenever the default database changes. If there is no default database, the variable has the same value as `character_set_server`.

The global `character_set_database` and `collation_database` system variables are deprecated and will be removed in a future version of MySQL.

Assigning a value to the session `character_set_database` and `collation_database` system variables is deprecated and assignments produce a warning. The session variables will become read only in a future version of MySQL and assignments will produce an error. It will remain possible to access the session variables to determine the database character set and collation for the default database.

- `character_set_filesystem`

Command-Line Format	`--character-set-filesystem=name`	
System Variable	**Name**	`character_set_filesystem`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`binary`

The file system character set. This variable is used to interpret string literals that refer to file names, such as in the `LOAD DATA INFILE` and `SELECT ... INTO OUTFILE` statements and the `LOAD_FILE()` function. Such file names are converted from `character_set_client` to `character_set_filesystem` before the file opening attempt occurs. The default value is `binary`, which means that no conversion occurs. For systems on which multibyte file names are

permitted, a different value may be more appropriate. For example, if the system represents file names using UTF-8, set `character_set_filesystem` to `'utf8'`.

- `character_set_results`

System Variable	Name	character_set_results
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string

The character set used for returning query results such as result sets or error messages to the client.

- `character_set_server`

Command-Line Format	--character-set-server	
System Variable	Name	character_set_server
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	Type	string
	Default	latin1
Permitted Values (>= 8.0.1)	Type	string
	Default	utf8mb4

The server's default character set.

- `character_set_system`

System Variable	Name	character_set_system
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	utf8

The character set used by the server for storing identifiers. The value is always `utf8`.

- `character_sets_dir`

Command-Line Format	--character-sets-dir=dir_name	
System Variable	Name	character_sets_dir
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The directory where character sets are installed.

- `check_proxy_users`

Command-Line Format	`--check-proxy-users=[={OFF|ON}]`	
System Variable	Name	`check_proxy_users`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

This variable controls whether the server performs proxy user mapping for authentication plugins that request it. With `check_proxy_users` enabled, it may also be necessary to enable plugin-specific system variables to take advantage of server proxy user mapping support:

- For the `mysql_native_password` plugin, enable `mysql_native_password_proxy_users`.
- For the `sha256_password` plugin, enable `sha256_password_proxy_users`.

For information about user proxying, see Section 6.3.11, “Proxy Users”.

- `collation_connection`

System Variable	Name	`collation_connection`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string

The collation of the connection character set.

- `collation_database`

System Variable	Name	`collation_database`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	Type	string
	Default	`latin1_swedish_ci`
Permitted Values (>= 8.0.1)	Type	string
	Default	`utf8mb4_0900_ai_ci`
Footnote	This option is dynamic, but only the server should set this information. You should not set the value of this variable manually.	

The collation used by the default database. The server sets this variable whenever the default database changes. If there is no default database, the variable has the same value as `collation_server`.

The global `character_set_database` and `collation_database` system variables are deprecated and will be removed in a future version of MySQL.

Assigning a value to the session `character_set_database` and `collation_database` system variables is deprecated and assignments produce a warning. The session variables will become read only in a future version of MySQL and assignments will produce an error. It will remain possible to access the session variables to determine the database character set and collation for the default database.

- `collation_server`

Command-Line Format	`--collation-server`	
System Variable	**Name**	`collation_server`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (8.0.0)	**Type**	string
	Default	`latin1_swedish_ci`
Permitted Values (>= 8.0.1)	**Type**	string
	Default	`utf8mb4_0900_ai_ci`

The server's default collation.

- `completion_type`

Command-Line Format	`--completion-type=#`	
System Variable	**Name**	`completion_type`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`NO_CHAIN`
	Valid Values	`NO_CHAIN`
		`CHAIN`
		`RELEASE`
		`0`
		`1`
		`2`

The transaction completion type. This variable can take the values shown in the following table. The variable can be assigned using either the name values or corresponding integer values.

Value	Description
`NO_CHAIN` (or 0)	`COMMIT` and `ROLLBACK` are unaffected. This is the default value.
`CHAIN` (or 1)	`COMMIT` and `ROLLBACK` are equivalent to `COMMIT AND CHAIN` and `ROLLBACK AND CHAIN`, respectively. (A new transaction starts immediately with the same isolation level as the just-terminated transaction.)
`RELEASE` (or 2)	`COMMIT` and `ROLLBACK` are equivalent to `COMMIT RELEASE` and `ROLLBACK RELEASE`, respectively. (The server disconnects after terminating the transaction.)

`completion_type` affects transactions that begin with `START TRANSACTION` or `BEGIN` and end with `COMMIT` or `ROLLBACK`. It does not apply to implicit commits resulting from execution of the statements listed in Section 13.3.3, “Statements That Cause an Implicit Commit”. It also does not apply for `XA COMMIT`, `XA ROLLBACK`, or when `autocommit=1`.

- `concurrent_insert`

Command-Line Format	`--concurrent-insert[=#]`	
System Variable	Name	`concurrent_insert`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`AUTO`
	Valid Values	`NEVER`
		`AUTO`
		`ALWAYS`
		`0`
		`1`
		`2`

If `AUTO` (the default), MySQL permits `INSERT` and `SELECT` statements to run concurrently for `MyISAM` tables that have no free blocks in the middle of the data file. If you start `mysqld` with `--skip-new`, this variable is set to `NEVER`.

This variable can take the values shown in the following table. The variable can be assigned using either the name values or corresponding integer values.

Value	Description
`NEVER` (or 0)	Disables concurrent inserts
`AUTO` (or 1)	(Default) Enables concurrent insert for `MyISAM` tables that do not have holes
`ALWAYS` (or 2)	Enables concurrent inserts for all `MyISAM` tables, even those that have holes. For a table with a hole, new rows are inserted at the end of the table if it is in use by another thread. Otherwise, MySQL acquires a normal write lock and inserts the row into the hole.

See also Section 8.11.3, “Concurrent Inserts”.

- `connect_timeout`

Command-Line Format	`--connect-timeout=#`	
System Variable	Name	`connect_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer

	Default	`10`
	Min Value	`2`
	Max Value	`31536000`

The number of seconds that the `mysqld` server waits for a connect packet before responding with `Bad handshake`. The default value is 10 seconds.

Increasing the `connect_timeout` value might help if clients frequently encounter errors of the form `Lost connection to MySQL server at 'XXX', system error: errno`.

- `core_file`

System Variable	**Name**	`core_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Whether to write a core file if the server crashes. This variable is set by the `--core-file` option.

- `datadir`

Command-Line Format	`--datadir=dir_name`	
System Variable	**Name**	`datadir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The path to the MySQL server data directory. Relative paths are resolved with respect to the current directory. If the server will be started automatically (that is, in contexts for which you cannot assume what the current directory will be), it is best to specify the `datadir` value as an absolute path.

- `date_format`

This variable is unused. It is deprecated and will be removed in a future MySQL release.

- `datetime_format`

This variable is unused. It is deprecated and will be removed in a future MySQL release.

- `debug`

Command-Line Format	`--debug[=debug_options]`	
System Variable	**Name**	`debug`
	Variable Scope	Global, Session
	Dynamic Variable	Yes

Permitted Values (Unix)	**Type**	string
	Default	`d:t:i:o,/tmp/mysqld.trace`
Permitted Values (Windows)	**Type**	string
	Default	`d:t:i:O,\mysqld.trace`

This variable indicates the current debugging settings. It is available only for servers built with debugging support. The initial value comes from the value of instances of the `--debug` option given at server startup. The global and session values may be set at runtime. The `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege is required, even for the session value.

Assigning a value that begins with `+` or `-` cause the value to added to or subtracted from the current value:

```
mysql> SET debug = 'T';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| T       |
+---------+

mysql> SET debug = '+P';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| P:T     |
+---------+

mysql> SET debug = '-P';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| T       |
+---------+
```

For more information, see Section 28.5.3, “The DBUG Package”.

- `debug_sync`

System Variable	**Name**	`debug_sync`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string

This variable is the user interface to the Debug Sync facility. Use of Debug Sync requires that MySQL be configured with the `-DENABLE_DEBUG_SYNC=1 CMake` option (see Section 2.8.4, “MySQL Source-Configuration Options”). If Debug Sync is not compiled in, this system variable is not available.

The global variable value is read only and indicates whether the facility is enabled. By default, Debug Sync is disabled and the value of `debug_sync` is `OFF`. If the server is started with `--debug-sync-timeout=N`, where *`N`* is a timeout value greater than 0, Debug Sync is enabled and the value of `debug_sync` is `ON - current signal` followed by the signal name. Also, *`N`* becomes the default timeout for individual synchronization points.

The session value can be read by any user and will have the same value as the global variable. The session value can be set by users that have the SYSTEM_VARIABLES_ADMIN or SUPER privilege to control synchronization points.

For a description of the Debug Sync facility and how to use synchronization points, see MySQL Internals: Test Synchronization.

- default_authentication_plugin

Command-Line Format	--default-authentication-plugin=plugin_name	
System Variable	Name	default_authentication_plugin
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	enumeration
	Default	mysql_native_password
	Valid Values	mysql_native_password
		sha256_password

The default authentication plugin. Permitted values are mysql_native_password (use MySQL native passwords; this is the default) and sha256_password (use SHA-256 passwords). For more information about these plugins, see Section 6.5.1.1, “Native Pluggable Authentication”, and Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

Note

If you use this variable to change the default authentication plugin to a value other than mysql_native_password, clients older than MySQL 5.5.6 will no longer be able to connect because they will not understand the resulting change to the authentication protocol.

The value of default_authentication_plugin affects these aspects of server operation:

- It determines which authentication plugin the server assigns to new accounts created by CREATE USER and GRANT statements that do not name a plugin explicitly with an IDENTIFIED WITH clause.
- It sets the old_passwords system variable at startup to the value that is consistent with the password hashing method required by the default plugin. The old_passwords value affects hashing of passwords specified in the IDENTIFIED BY clause of CREATE USER and GRANT, and passwords specified as the argument to the PASSWORD() function.
- For an account created with either of the following statements, the server associates the account with the default authentication plugin and assigns the account the given password, hashed according to the value of old_passwords.

```
CREATE USER ... IDENTIFIED BY 'cleartext password';
GRANT ...  IDENTIFIED BY 'cleartext password';
```

- For an account created with either of the following statements, the statement fails if the password hash is not encrypted using the hash format required by the default authentication plugin. Otherwise, the server associates the account with the default authentication plugin and assigns the account the given password hash.

```
CREATE USER ... IDENTIFIED BY PASSWORD 'encrypted password';
GRANT ...  IDENTIFIED BY PASSWORD 'encrypted password';
```

- `default_password_lifetime`

Command-Line Format	`--default-password-lifetime=#`	
System Variable	Name	`default_password_lifetime`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`65535`

This variable defines the global automatic password expiration policy. It applies to accounts that use MySQL built-in authentication methods (accounts that use an authentication plugin of `mysql_native_password` or `sha256_password`).

The default `default_password_lifetime` value is 0, which disables automatic password expiration. If the value of `default_password_lifetime` is a positive integer *`N`*, it indicates the permitted password lifetime; passwords must be changed every *`N`* days.

The global password expiration policy can be overridden as desired for individual accounts using the `ALTER USER` statement. See Section 6.3.8, “Password Expiration Policy”.

- `default_storage_engine`

Command-Line Format	`--default-storage-engine=name`	
System Variable	Name	`default_storage_engine`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`InnoDB`

The default storage engine. This variable sets the storage engine for permanent tables only. To set the storage engine for `TEMPORARY` tables, set the `default_tmp_storage_engine` system variable.

To see which storage engines are available and enabled, use the `SHOW ENGINES` statement or query the `INFORMATION_SCHEMA ENGINES` table.

If you disable the default storage engine at server startup, you must set the default engine for both permanent and `TEMPORARY` tables to a different engine or the server will not start.

- `default_tmp_storage_engine`

Command-Line Format	`--default-tmp-storage-engine=name`	
System Variable	Name	`default_tmp_storage_engine`

	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`InnoDB`

The default storage engine for `TEMPORARY` tables (created with `CREATE TEMPORARY TABLE`). To set the storage engine for permanent tables, set the `default_storage_engine` system variable. Also see the discussion of that variable regarding possible values.

If you disable the default storage engine at server startup, you must set the default engine for both permanent and `TEMPORARY` tables to a different engine or the server will not start.

- `default_week_format`

Command-Line Format	`--default-week-format=#`	
System Variable	**Name**	`default_week_format`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`7`

The default mode value to use for the `WEEK()` function. See Section 12.7, “Date and Time Functions”.

- `delay_key_write`

Command-Line Format	`--delay-key-write[=name]`	
System Variable	**Name**	`delay_key_write`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`ON`
	Valid Values	`ON`
		`OFF`
		`ALL`

This option applies only to `MyISAM` tables. It can have one of the following values to affect handling of the `DELAY_KEY_WRITE` table option that can be used in `CREATE TABLE` statements.

Option	Description
`OFF`	`DELAY_KEY_WRITE` is ignored.

Option	Description
ON	MySQL honors any DELAY_KEY_WRITE option specified in CREATE TABLE statements. This is the default value.
ALL	All new opened tables are treated as if they were created with the DELAY_KEY_WRITE option enabled.

If DELAY_KEY_WRITE is enabled for a table, the key buffer is not flushed for the table on every index update, but only when the table is closed. This speeds up writes on keys a lot, but if you use this feature, you should add automatic checking of all MyISAM tables by starting the server with the --myisam-recover-options option (for example, --myisam-recover-options=BACKUP,FORCE). See Section 5.1.4, "Server Command Options", and Section 16.2.1, "MyISAM Startup Options".

Warning

If you enable external locking with --external-locking, there is no protection against index corruption for tables that use delayed key writes.

- delayed_insert_limit

Deprecated	5.6.7	
Command-Line Format	--delayed-insert-limit=#	
System Variable	**Name**	delayed_insert_limit
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	100
	Min Value	1
	Max Value	4294967295
Permitted Values (64-bit platforms)	**Type**	integer
	Default	100
	Min Value	1
	Max Value	18446744073709551615

This system variable is deprecated (because DELAYED inserts are not supported), and will be removed in a future release.

- delayed_insert_timeout

Deprecated	5.6.7	
Command-Line Format	--delayed-insert-timeout=#	
System Variable	**Name**	delayed_insert_timeout
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`300`

This system variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `delayed_queue_size`

Deprecated	5.6.7	
Command-Line Format	`--delayed-queue-size=#`	
System Variable	**Name**	`delayed_queue_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`1000`
	Min Value	`1`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`1000`
	Min Value	`1`
	Max Value	`18446744073709551615`

This system variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `disabled_storage_engines`

Command-Line Format	`--disabled-storage-engines=engine[,engine]...`	
System Variable	**Name**	`disabled_storage_engines`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`empty string`

This variable indicates which storage engines cannot be used to create tables or tablespaces. For example, to prevent new `MyISAM` or `FEDERATED` tables from being created, start the server with these lines in the server option file:

```
[mysqld]
disabled_storage_engines="MyISAM,FEDERATED"
```

By default, `disabled_storage_engines` is empty (no engines disabled), but it can be set to a comma-separated list of one or more engines (not case sensitive). Any engine named in the value cannot be used to create tables or tablespaces with `CREATE TABLE` or `CREATE TABLESPACE`, and cannot be used with `ALTER TABLE ... ENGINE` or `ALTER TABLESPACE ... ENGINE` to change the storage engine of existing tables or tablespaces. Attempts to do so result in an `ER_DISABLED_STORAGE_ENGINE` error.

`disabled_storage_engines` does not restrict other DDL statements for existing tables, such as `CREATE INDEX`, `TRUNCATE TABLE`, `ANALYZE TABLE`, `DROP TABLE`, or `DROP TABLESPACE`. This permits a smooth transition so that existing tables or tablespaces that use a disabled engine can be migrated to a permitted engine by means such as `ALTER TABLE ... ENGINE` *`permitted_engine`*.

It is permitted to set the `default_storage_engine` or `default_tmp_storage_engine` system variable to a storage engine that is disabled. This could cause applications to behave erratically or fail, although that might be a useful technique in a development environment for identifying applications that use disabled engines, so that they can be modified.

`disabled_storage_engines` is disabled and has no effect if the server is started with any of these options: `--initialize`, `--initialize-insecure`, `--skip-grant-tables`.

- `disconnect_on_expired_password`

Command-Line Format	`--disconnect-on-expired-password[=#]`	
System Variable	**Name**	`disconnect_on_expired_password`
	Variable Scope	Session
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`ON`

This variable controls how the server handles clients with expired passwords:

- If the client indicates that it can handle expires passwords, the value of `disconnect_on_expired_password` is irrelevant. The server permits the client to connect but puts it in sandbox mode.
- If the client does not indicate that it can handle expires passwords, the server handles the client according to the value of `disconnect_on_expired_password`:
  - If `disconnect_on_expired_password`: is enabled, the server disconnects the client.
  - If `disconnect_on_expired_password`: is disabled, the server permits the client to connect but puts it in sandbox mode.

For more information about the interaction of client and server settings relating to expired-password handling, see Section 6.3.9, “Password Expiration and Sandbox Mode”.

- `div_precision_increment`

Command-Line Format	`--div-precision-increment=#`	
System Variable	**Name**	`div_precision_increment`
	Variable Scope	Global, Session

	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`4`
	Min Value	`0`
	Max Value	`30`

This variable indicates the number of digits by which to increase the scale of the result of division operations performed with the `/` operator. The default value is 4. The minimum and maximum values are 0 and 30, respectively. The following example illustrates the effect of increasing the default value.

```
mysql> SELECT 1/7;
+--------+
| 1/7    |
+--------+
| 0.1429 |
+--------+
mysql> SET div_precision_increment = 12;
mysql> SELECT 1/7;
+----------------+
| 1/7            |
+----------------+
| 0.142857142857 |
+----------------+
```

- `end_markers_in_json`

System Variable	**Name**	`end_markers_in_json`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Whether optimizer JSON output should add end markers. See MySQL Internals: The end_markers_in_json System Variable.

- `eq_range_index_dive_limit`

System Variable	**Name**	`eq_range_index_dive_limit`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`200`
	Min Value	`0`
	Max Value	`4294967295`

This variable indicates the number of equality ranges in an equality comparison condition when the optimizer should switch from using index dives to index statistics in estimating the number of qualifying rows. It applies to evaluation of expressions that have either of these equivalent forms, where the optimizer uses a nonunique index to look up *`col_name`* values:

```
col_name IN(val1, ..., valN)
col_name = val1 OR ... OR col_name = valN
```

In both cases, the expression contains *`N`* equality ranges. The optimizer can make row estimates using index dives or index statistics. If `eq_range_index_dive_limit` is greater than 0, the optimizer uses existing index statistics instead of index dives if there are `eq_range_index_dive_limit` or more equality ranges. Thus, to permit use of index dives for up to *`N`* equality ranges, set `eq_range_index_dive_limit` to *`N`* + 1. To disable use of index statistics and always use index dives regardless of *`N`*, set `eq_range_index_dive_limit` to 0.

For more information, see Equality Range Optimization of Many-Valued Comparisons.

To update table index statistics for best estimates, use `ANALYZE TABLE`.

- `error_count`

  The number of errors that resulted from the last statement that generated messages. This variable is read only. See Section 13.7.5.17, “SHOW ERRORS Syntax”.

- `event_scheduler`

Command-Line Format	`--event-scheduler[=value]`	
System Variable	**Name**	`event_scheduler`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`OFF`
	Valid Values	`ON`
		`OFF`
		`DISABLED`

This variable indicates the status of the Event Scheduler; possible values are `ON`, `OFF`, and `DISABLED`, with the default being `OFF`. This variable and its effects on the Event Scheduler's operation are discussed in greater detail in the Overview section of the Events chapter.

- `expire_logs_days`

Command-Line Format	`--expire-logs-days=#`	
System Variable	**Name**	`expire_logs_days`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`

	Min Value	0
	Max Value	99
Permitted Values (>= 8.0.2)	**Type**	integer
	Default	30
	Min Value	0
	Max Value	99

The number of days before automatic removal of binary log files. The default is 30, which means binary log files are removed after 30 days. Setting this variable to 0 means “no automatic removal”. Possible removals happen at startup and when the binary log is flushed. Log flushing occurs as indicated in Section 5.4, “MySQL Server Logs”.

This variable sets the binary log expiration period in whole days. To set this interval to less than one day, use `binlog_expire_logs_seconds`. The effect of the two variables are cumulative; for example, if `expire_logs_days` = 1 and `binlog_expire_logs_seconds` = 43200, then the binary log is purged every 1.5 days.

In MySQL 8.0, both `binlog_expire_logs_seconds` and `expire_logs_days` must be set to 0 to disable automatic purging of the binary log.

To remove binary log files manually, use the `PURGE BINARY LOGS` statement. See Section 13.4.1.1, “PURGE BINARY LOGS Syntax”.

- `explicit_defaults_for_timestamp`

Deprecated	5.6.6	
Command-Line Format	`--explicit-defaults-for-timestamp=#`	
System Variable	**Name**	`explicit_defaults_for_timestamp`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (<= 8.0.1)	**Type**	boolean
	Default	`OFF`
Permitted Values (>= 8.0.2)	**Type**	boolean
	Default	`ON`

This system variable determines whether the server enables certain nonstandard behaviors for default values and `NULL`-value handling in `TIMESTAMP` columns. By default, `explicit_defaults_for_timestamp` is enabled, which disables the nonstandard behaviors. Disabling `explicit_defaults_for_timestamp` results in a warning. (Before MySQL 8.0.2, `explicit_defaults_for_timestamp` is disabled, which enables the nonstandard behaviors.)

If `explicit_defaults_for_timestamp` is disabled, the server enables the nonstandard behaviors and handles `TIMESTAMP` columns as follows:

- `TIMESTAMP` columns not explicitly declared with the `NULL` attribute are automatically declared with the `NOT NULL` attribute. Assigning such a column a value of `NULL` is permitted and sets the column to the current timestamp.

- The first `TIMESTAMP` column in a table, if not explicitly declared with the `NULL` attribute or an explicit `DEFAULT` or `ON UPDATE` attribute, is automatically declared with the `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP` attributes.

- `TIMESTAMP` columns following the first one, if not explicitly declared with the `NULL` attribute or an explicit `DEFAULT` attribute, are automatically declared as `DEFAULT '0000-00-00 00:00:00'` (the "zero" timestamp). For inserted rows that specify no explicit value for such a column, the column is assigned `'0000-00-00 00:00:00'` and no warning occurs.

  Depending on whether strict SQL mode or the `NO_ZERO_DATE` SQL mode is enabled, a default value of `'0000-00-00 00:00:00'` may be invalid. Be aware that the `TRADITIONAL` SQL mode includes strict mode and `NO_ZERO_DATE`. See Section 5.1.8, "Server SQL Modes".

The nonstandard behaviors just described are deprecated and will be removed in a future MySQL release.

If `explicit_defaults_for_timestamp` is enabled, the server disables the nonstandard behaviors and handles `TIMESTAMP` columns as follows:

- It is not possible to assign a `TIMESTAMP` column a value of `NULL` to set it to the current timestamp. To assign the current timestamp, set the column to `CURRENT_TIMESTAMP` or a synonym such as `NOW()`.

- `TIMESTAMP` columns not explicitly declared with the `NOT NULL` attribute are automatically declared with the `NULL` attribute and permit `NULL` values. Assigning such a column a value of `NULL` sets it to `NULL`, not the current timestamp.

- `TIMESTAMP` columns declared with the `NOT NULL` attribute do not permit `NULL` values. For inserts that specify `NULL` for such a column, the result is either an error for a single-row insert or if strict SQL mode is enabled, or `'0000-00-00 00:00:00'` is inserted for multiple-row inserts with strict SQL mode disabled. In no case does assigning the column a value of `NULL` set it to the current timestamp.

- `TIMESTAMP` columns explicitly declared with the `NOT NULL` attribute and without an explicit `DEFAULT` attribute are treated as having no default value. For inserted rows that specify no explicit value for such a column, the result depends on the SQL mode. If strict SQL mode is enabled, an error occurs. If strict SQL mode is not enabled, the column is declared with the implicit default of `'0000-00-00 00:00:00'` and a warning occurs. This is similar to how MySQL treats other temporal types such as `DATETIME`.

- No `TIMESTAMP` column is automatically declared with the `DEFAULT CURRENT_TIMESTAMP` or `ON UPDATE CURRENT_TIMESTAMP` attributes. Those attributes must be explicitly specified.

- The first `TIMESTAMP` column in a table is not handled differently from `TIMESTAMP` columns following the first one.

If `explicit_defaults_for_timestamp` is disabled at server startup, this warning appears in the error log:

```
[Warning] TIMESTAMP with implicit DEFAULT value is deprecated.
Please use --explicit_defaults_for_timestamp server option (see
documentation for more details).
```

As indicated by the warning, to disable the deprecated nonstandard behaviors, enable the `explicit_defaults_for_timestamp` system variable at server startup.

Note

`explicit_defaults_for_timestamp` is itself deprecated because its only purpose is to permit control over deprecated `TIMESTAMP` behaviors that are to be removed in a future MySQL release. When removal of those behaviors occurs, `explicit_defaults_for_timestamp` will have no purpose and will be removed as well.

For additional information, see Section 11.3.5, "Automatic Initialization and Updating for TIMESTAMP and DATETIME".

- `external_user`

System Variable	Name	external_user
	Variable Scope	Session
	Dynamic Variable	No
Permitted Values	Type	string

The external user name used during the authentication process, as set by the plugin used to authenticate the client. With native (built-in) MySQL authentication, or if the plugin does not set the value, this variable is `NULL`. See Section 6.3.11, "Proxy Users".

- `flush`

Command-Line Format	--flush	
System Variable	Name	flush
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

If `ON`, the server flushes (synchronizes) all changes to disk after each SQL statement. Normally, MySQL does a write of all changes to disk only after each SQL statement and lets the operating system handle the synchronizing to disk. See Section B.5.3.3, "What to Do If MySQL Keeps Crashing". This variable is set to `ON` if you start `mysqld` with the `--flush` option.

- `flush_time`

Command-Line Format	--flush-time=#	
System Variable	Name	flush_time
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	0
	Min	0
	Value	

Permitted Values (Windows)	**Type**	integer
	Default	`0`
	Min Value	`0`

If this is set to a nonzero value, all tables are closed every `flush_time` seconds to free up resources and synchronize unflushed data to disk. This option is best used only on systems with minimal resources.

- `foreign_key_checks`

System Variable	**Name**	`foreign_key_checks`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`1`

If set to 1 (the default), foreign key constraints for `InnoDB` tables are checked. If set to 0, foreign key constraints are ignored, with a couple of exceptions. When re-creating a table that was dropped, an error is returned if the table definition does not conform to the foreign key constraints referencing the table. Likewise, an `ALTER TABLE` operation returns an error if a foreign key definition is incorrectly formed. For more information, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”.

Typically you leave this setting enabled during normal operation, to enforce referential integrity. Disabling foreign key checking can be useful for reloading `InnoDB` tables in an order different from that required by their parent/child relationships. See Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

Setting `foreign_key_checks` to 0 also affects data definition statements: `DROP SCHEMA` drops a schema even if it contains tables that have foreign keys that are referred to by tables outside the schema, and `DROP TABLE` drops tables that have foreign keys that are referred to by other tables.

Note

Setting `foreign_key_checks` to 1 does not trigger a scan of the existing table data. Therefore, rows added to the table while `foreign_key_checks = 0` will not be verified for consistency.

Dropping an index required by a foreign key constraint is not permitted, even with `foreign_key_checks=0`. The foreign key constraint must be removed before dropping the index.

- `ft_boolean_syntax`

Command-Line Format	`--ft-boolean-syntax=name`	
System Variable	**Name**	`ft_boolean_syntax`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`+ -><()~*:""&|`

The list of operators supported by boolean full-text searches performed using `IN BOOLEAN MODE`. See Section 12.9.2, “Boolean Full-Text Searches”.

The default variable value is `'+ -><()~*:""&|'`. The rules for changing the value are as follows:

- Operator function is determined by position within the string.
- The replacement value must be 14 characters.
- Each character must be an ASCII nonalphanumeric character.
- Either the first or second character must be a space.
- No duplicates are permitted except the phrase quoting operators in positions 11 and 12. These two characters are not required to be the same, but they are the only two that may be.
- Positions 10, 13, and 14 (which by default are set to `:`, `&`, and `|`) are reserved for future extensions.

- `ft_max_word_len`

Command-Line Format	`--ft-max-word-len=#`	
System Variable	Name	`ft_max_word_len`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Min Value	`10`

The maximum length of the word to be included in a `MyISAM FULLTEXT` index.

Note

`FULLTEXT` indexes on `MyISAM` tables must be rebuilt after changing this variable. Use `REPAIR TABLE tbl_name QUICK`.

- `ft_min_word_len`

Command-Line Format	`--ft-min-word-len=#`	
System Variable	Name	`ft_min_word_len`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`4`
	Min Value	`1`

The minimum length of the word to be included in a `MyISAM FULLTEXT` index.

Note

`FULLTEXT` indexes on `MyISAM` tables must be rebuilt after changing this variable. Use `REPAIR TABLE tbl_name QUICK`.

- `ft_query_expansion_limit`

Command-Line Format	`--ft-query-expansion-limit=#`	
System Variable	Name	`ft_query_expansion_limit`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`20`
	Min Value	`0`
	Max Value	`1000`

The number of top matches to use for full-text searches performed using `WITH QUERY EXPANSION`.

- `ft_stopword_file`

Command-Line Format	`--ft-stopword-file=file_name`	
System Variable	Name	`ft_stopword_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The file from which to read the list of stopwords for full-text searches on `MyISAM` tables. The server looks for the file in the data directory unless an absolute path name is given to specify a different directory. All the words from the file are used; comments are *not* honored. By default, a built-in list of stopwords is used (as defined in the `storage/myisam/ft_static.c` file). Setting this variable to the empty string (`''`) disables stopword filtering. See also Section 12.9.4, “Full-Text Stopwords”.

Note

`FULLTEXT` indexes on `MyISAM` tables must be rebuilt after changing this variable or the contents of the stopword file. Use `REPAIR TABLE tbl_name QUICK`.

- `general_log`

Command-Line Format	`--general-log`	
System Variable	Name	`general_log`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	Type	boolean
	Default	OFF

Whether the general query log is enabled. The value can be 0 (or OFF) to disable the log or 1 (or ON) to enable the log. The default value depends on whether the --general_log option is given. The destination for log output is controlled by the log_output system variable; if that value is NONE, no log entries are written even if the log is enabled.

- general_log_file

Command-Line Format	--general-log-file=file_name	
System Variable	Name	general_log_file
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	file name
	Default	host_name.log

The name of the general query log file. The default value is *host_name*.log, but the initial value can be changed with the --general_log_file option.

- group_concat_max_len

Command-Line Format	--group-concat-max-len=#	
System Variable	Name	group_concat_max_len
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	1024
	Min Value	4
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	1024
	Min Value	4
	Max Value	18446744073709551615

The maximum permitted result length in bytes for the GROUP_CONCAT() function. The default is 1024.

- have_compress

YES if the zlib compression library is available to the server, NO if not. If not, the COMPRESS() and UNCOMPRESS() functions cannot be used.

- have_crypt

YES if the crypt() system call is available to the server, NO if not. If not, the ENCRYPT() function cannot be used.

- have_dynamic_loading

  YES if mysqld supports dynamic loading of plugins, NO if not. If the value is NO, you cannot use options such as --plugin-load to load plugins at server startup, or the INSTALL PLUGIN statement to load plugins at runtime.

- have_geometry

  YES if the server supports spatial data types, NO if not.

- have_openssl

  This variable is an alias for have_ssl.

- have_profiling

  YES if statement profiling capability is present, NO if not. If present, the profiling system variable controls whether this capability is enabled or disabled. See Section 13.7.5.31, “SHOW PROFILES Syntax”.

  This variable is deprecated and will be removed in a future MySQL release.

- have_query_cache

  YES if mysqld supports the query cache, NO if not.

- have_rtree_keys

  YES if RTREE indexes are available, NO if not. (These are used for spatial indexes in MyISAM tables.)

- have_ssl

  YES if mysqld supports SSL connections, NO if not. DISABLED indicates that the server was compiled with SSL support, but was not started with the appropriate --ssl-*xxx* options. For more information, see Section 6.4.2, “Building MySQL with Support for Secure Connections”.

- have_statement_timeout

System Variable	Name	have_statement_timeout
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean

  Whether the statement execution timeout feature is available (see Statement Execution Time Optimizer Hints). The value can be NO if the background thread used by this feature could not be initialized.

- have_symlink

  YES if symbolic link support is enabled, NO if not. This is required on Unix for support of the DATA DIRECTORY and INDEX DIRECTORY table options. If the server is started with the --skip-symbolic-links option, the value is DISABLED.

  This variable has no meaning on Windows.

Note

As of MySQL 8.0.2, symbolic link support, along with the the `--symbolic-links` option that controls it, is deprecated and will be removed in a future version of MySQL. In addition, the option is disabled by default. The related `have_symlink` system variable also is deprecated and will be removed in a future version of MySQL.

- `histogram_generation_max_mem_size`

Introduced	8.0.2	
Command-Line Format	`--histogram-generation-max-mem-size=#`	
System Variable	**Name**	`histogram_generation_max_mem_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`20000000`
	Min Value	`1000000`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`20000000`
	Min Value	`1000000`
	Max Value	`18446744073709551615`

The maximum amount of memory available for generating histogram statistics. See Section 8.9.6, “Optimizer Statistics”, and Section 13.7.2.1, “ANALYZE TABLE Syntax”.

The global and session values may be set at runtime. The `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege is required, even for the session value.

- `host_cache_size`

System Variable	**Name**	`host_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`
	Min Value	`0`
	Max Value	`65536`

The size of the internal host cache (see Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”). Setting the size to 0 disables the host cache. Changing the cache size at runtime implicitly causes a `FLUSH HOSTS` operation to clear the host cache and truncate the `host_cache` table.

The default value is 128, plus 1 for a value of `max_connections` up to 500, plus 1 for every increment of 20 over 500 in the `max_connections` value, capped to a limit of 2000.

Use of `--skip-host-cache` is similar to setting the `host_cache_size` system variable to 0, but `host_cache_size` is more flexible because it can also be used to resize, enable, or disable the host cache at runtime, not just at server startup.

If you start the server with `--skip-host-cache`, that does not prevent changes to the value of `host_cache_size`, but such changes have no effect and the cache is not re-enabled even if `host_cache_size` is set larger than 0.

- `hostname`

System Variable	Name	hostname
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The server sets this variable to the server host name at startup.

- `identity`

This variable is a synonym for the `last_insert_id` variable. It exists for compatibility with other database systems. You can read its value with `SELECT @@identity`, and set it using `SET identity`.

- `init_connect`

Command-Line Format	--init-connect=name	
System Variable	Name	init_connect
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

A string to be executed by the server for each client that connects. The string consists of one or more SQL statements, separated by semicolon characters. For example, each client session begins by default with autocommit mode enabled. For older servers (before MySQL 5.5.8), there is no global `autocommit` system variable to specify that autocommit should be disabled by default, but as a workaround `init_connect` can be used to achieve the same effect:

```
SET GLOBAL init_connect='SET autocommit=0';
```

The `init_connect` variable can also be set on the command line or in an option file. To set the variable as just shown using an option file, include these lines:

```
[mysqld]
init_connect='SET autocommit=0'
```

For users that have the CONNECTION_ADMIN or SUPER privilege, the content of init_connect is not executed. This is done so that an erroneous value for init_connect does not prevent all clients from connecting. For example, the value might contain a statement that has a syntax error, thus causing client connections to fail. Not executing init_connect for users that have the CONNECTION_ADMIN or SUPER privilege enables them to open a connection and fix the init_connect value.

The server discards any result sets produced by statements in the value of of init_connect.

- information_schema_stats

Command-Line Format	--information-schema-stats=value	
System Variable	Name	information_schema_stats
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	CACHED
	Valid Values	CACHED
		LATEST

Some INFORMATION_SCHEMA tables contain columns that provide table statistics:

```
STATISTICS.CARDINALITY
TABLES.AUTO_INCREMENT
TABLES.AVG_ROW_LENGTH
TABLES.CHECKSUM
TABLES.CHECK_TIME
TABLES.CREATE_TIME
TABLES.DATA_FREE
TABLES.DATA_LENGTH
TABLES.INDEX_LENGTH
TABLES.MAX_DATA_LENGTH
TABLES.TABLE_ROWS
TABLES.UPDATE_TIME
```

Those columns represent dynamic table metadata; that is, information that changes as table contents change.

The information_schema_stats system variable determines the source from which the server retrieves table statistics:

- If the value is CACHED (the default), the server retrieves statistics from INFORMATION_SCHEMA tables that hold cached values (STATISTICS and TABLES). When the server starts, the cached statistics are NULL. To update the cached values for a given table, use ANALYZE TABLE.
- If the value is LATEST, the server reads statistics directly from storage engines.

Note

If the innodb_read_only system variable is enabled, ANALYZE TABLE may fail because it cannot update statistics tables in the data dictionary, which use InnoDB. For ANALYZE TABLE operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a MyISAM table). To obtain the updated distribution statistics, set information_schema_stats=LATEST.

For additional information, see Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”.

- `init_file`

Command-Line Format	`--init-file=file_name`	
System Variable	Name	`init_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The name of the file specified with the `--init-file` option when you start the server. This should be a file containing SQL statements that you want the server to execute when it starts. Each statement must be on a single line and should not include comments. For more information, see the description of `--init-file`.

- `innodb_xxx`

`InnoDB` system variables are listed in Section 15.13, “InnoDB Startup Options and System Variables”. These variables control many aspects of storage, memory use, and I/O patterns for `InnoDB` tables, and are especially important now that `InnoDB` is the default storage engine.

- `insert_id`

The value to be used by the following `INSERT` or `ALTER TABLE` statement when inserting an `AUTO_INCREMENT` value. This is mainly used with the binary log.

- `interactive_timeout`

Command-Line Format	`--interactive-timeout=#`	
System Variable	Name	`interactive_timeout`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`28800`
	Min Value	`1`

The number of seconds the server waits for activity on an interactive connection before closing it. An interactive client is defined as a client that uses the `CLIENT_INTERACTIVE` option to `mysql_real_connect()`. See also `wait_timeout`.

- `internal_tmp_disk_storage_engine`

Command-Line Format	`--internal-tmp-disk-storage-engine=#`	
System Variable	Name	`internal_tmp_disk_storage_engine`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	enumeration
	Default	`INNODB`
	Valid Values	`MYISAM`
		`INNODB`

The storage engine for on-disk internal temporary tables (see Section 8.4.4, “Internal Temporary Table Use in MySQL”). Permitted values are `MYISAM` and `INNODB` (the default).

Under `internal_tmp_disk_storage_engine=INNODB`, queries that generate temporary tables that exceed `InnoDB` row or column limits will return `Row size too large` or `Too many columns` errors. The workaround is to set `internal_tmp_disk_storage_engine` to `MYISAM`.

- `internal_tmp_mem_storage_engine`

Introduced	8.0.2	
Command-Line Format	`--internal-tmp-mem-storage-engine=#`	
System Variable	**Name**	`internal_tmp_mem_storage_engine`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`TempTable`
	Valid Values	`TempTable`
		`MEMORY`

The storage engine for in-memory internal temporary tables (see Section 8.4.4, “Internal Temporary Table Use in MySQL”). Permitted values are `TempTable` (the default) and `MEMORY`.

- `join_buffer_size`

Command-Line Format	`--join-buffer-size=#`	
System Variable	**Name**	`join_buffer_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (Windows)	**Type**	integer
	Default	`262144`
	Min Value	`128`
	Max Value	`4294967295`
Permitted Values (Other, 32-bit platforms)	**Type**	integer
	Default	`262144`
	Min Value	`128`
	Max Value	`4294967295`

Permitted Values (Other, 64-bit platforms)	**Type**	integer
	Default	`262144`
	Min Value	`128`
	Max Value	`18446744073709547520`

The minimum size of the buffer that is used for plain index scans, range index scans, and joins that do not use indexes and thus perform full table scans. Normally, the best way to get fast joins is to add indexes. Increase the value of `join_buffer_size` to get a faster full join when adding indexes is not possible. One join buffer is allocated for each full join between two tables. For a complex join between several tables for which indexes are not used, multiple join buffers might be necessary.

Unless Batched Key Access (BKA) is used, there is no gain from setting the buffer larger than required to hold each matching row, and all joins allocate at least the minimum size, so use caution in setting this variable to a large value globally. It is better to keep the global setting small and change to a larger setting only in sessions that are doing large joins. Memory allocation time can cause substantial performance drops if the global size is larger than needed by most queries that use it.

When BKA is used, the value of `join_buffer_size` defines how large the batch of keys is in each request to the storage engine. The larger the buffer, the more sequential access will be to the right hand table of a join operation, which can significantly improve performance.

The default is 256KB. The maximum permissible setting for `join_buffer_size` is 4GB−1. Larger values are permitted for 64-bit platforms (except 64-bit Windows, for which large values are truncated to 4GB−1 with a warning).

For additional information about join buffering, see Section 8.2.1.6, “Nested-Loop Join Algorithms”. For information about Batched Key Access, see Section 8.2.1.11, “Block Nested-Loop and Batched Key Access Joins”.

- `keep_files_on_create`

Command-Line Format	`--keep-files-on-create=#`	
System Variable	**Name**	`keep_files_on_create`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

If a `MyISAM` table is created with no `DATA DIRECTORY` option, the `.MYD` file is created in the database directory. By default, if `MyISAM` finds an existing `.MYD` file in this case, it overwrites it. The same applies to `.MYI` files for tables created with no `INDEX DIRECTORY` option. To suppress this behavior, set the `keep_files_on_create` variable to `ON` (1), in which case `MyISAM` will not overwrite existing files and returns an error instead. The default value is `OFF` (0).

If a `MyISAM` table is created with a `DATA DIRECTORY` or `INDEX DIRECTORY` option and an existing `.MYD` or `.MYI` file is found, MyISAM always returns an error. It will not overwrite a file in the specified directory.

- `key_buffer_size`

Command-Line Format	`--key-buffer-size=#`

System Variable	Name	key_buffer_size
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	8388608
	Min Value	8
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	8388608
	Min Value	8
	Max Value	OS_PER_PROCESS_LIMIT

Index blocks for MyISAM tables are buffered and are shared by all threads. key_buffer_size is the size of the buffer used for index blocks. The key buffer is also known as the key cache.

The maximum permissible setting for key_buffer_size is 4GB−1 on 32-bit platforms. Larger values are permitted for 64-bit platforms. The effective maximum size might be less, depending on your available physical RAM and per-process RAM limits imposed by your operating system or hardware platform. The value of this variable indicates the amount of memory requested. Internally, the server allocates as much memory as possible up to this amount, but the actual allocation might be less.

You can increase the value to get better index handling for all reads and multiple writes; on a system whose primary function is to run MySQL using the MyISAM storage engine, 25% of the machine's total memory is an acceptable value for this variable. However, you should be aware that, if you make the value too large (for example, more than 50% of the machine's total memory), your system might start to page and become extremely slow. This is because MySQL relies on the operating system to perform file system caching for data reads, so you must leave some room for the file system cache. You should also consider the memory requirements of any other storage engines that you may be using in addition to MyISAM.

For even more speed when writing many rows at the same time, use LOCK TABLES. See Section 8.2.5.1, “Optimizing INSERT Statements”.

You can check the performance of the key buffer by issuing a SHOW STATUS statement and examining the Key_read_requests, Key_reads, Key_write_requests, and Key_writes status variables. (See Section 13.7.5, “SHOW Syntax”.) The Key_reads/Key_read_requests ratio should normally be less than 0.01. The Key_writes/Key_write_requests ratio is usually near 1 if you are using mostly updates and deletes, but might be much smaller if you tend to do updates that affect many rows at the same time or if you are using the DELAY_KEY_WRITE table option.

The fraction of the key buffer in use can be determined using key_buffer_size in conjunction with the Key_blocks_unused status variable and the buffer block size, which is available from the key_cache_block_size system variable:

```
1 - ((Key_blocks_unused * key_cache_block_size) / key_buffer_size)
```

This value is an approximation because some space in the key buffer is allocated internally for administrative structures. Factors that influence the amount of overhead for these structures include block size and pointer size. As block size increases, the percentage of the key buffer lost to overhead tends to decrease. Larger blocks results in a smaller number of read operations (because more keys are obtained per read), but conversely an increase in reads of keys that are not examined (if not all keys in a block are relevant to a query).

It is possible to create multiple `MyISAM` key caches. The size limit of 4GB applies to each cache individually, not as a group. See Section 8.10.2, “The MyISAM Key Cache”.

- `key_cache_age_threshold`

Command-Line Format	`--key-cache-age-threshold=#`	
System Variable	**Name**	`key_cache_age_threshold`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`300`
	Min Value	`100`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`300`
	Min Value	`100`
	Max Value	`18446744073709551615`

This value controls the demotion of buffers from the hot sublist of a key cache to the warm sublist. Lower values cause demotion to happen more quickly. The minimum value is 100. The default value is 300. See Section 8.10.2, “The MyISAM Key Cache”.

- `key_cache_block_size`

Command-Line Format	`--key-cache-block-size=#`	
System Variable	**Name**	`key_cache_block_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1024`
	Min Value	`512`
	Max Value	`16384`

The size in bytes of blocks in the key cache. The default value is 1024. See Section 8.10.2, “The MyISAM Key Cache”.

- `key_cache_division_limit`

Command-Line Format	`--key-cache-division-limit=#`	
System Variable	Name	`key_cache_division_limit`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`100`
	Min Value	`1`
	Max Value	`100`

The division point between the hot and warm sublists of the key cache buffer list. The value is the percentage of the buffer list to use for the warm sublist. Permissible values range from 1 to 100. The default value is 100. See Section 8.10.2, “The MyISAM Key Cache”.

- `large_files_support`

System Variable	Name	`large_files_support`
	Variable Scope	Global
	Dynamic Variable	No

Whether `mysqld` was compiled with options for large file support.

- `large_pages`

Command-Line Format	`--large-pages`	
System Variable	Name	`large_pages`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Linux	
Permitted Values (Linux)	Type	boolean
	Default	`FALSE`

Whether large page support is enabled (via the `--large-pages` option). See Section 8.12.3.2, “Enabling Large Page Support”.

- `large_page_size`

System Variable	Name	`large_page_size`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values (Linux)	**Type**	integer
	Default	`0`

If large page support is enabled, this shows the size of memory pages. Large memory pages are supported only on Linux; on other platforms, the value of this variable is always 0. See Section 8.12.3.2, “Enabling Large Page Support”.

- `last_insert_id`

  The value to be returned from `LAST_INSERT_ID()`. This is stored in the binary log when you use `LAST_INSERT_ID()` in a statement that updates a table. Setting this variable does not update the value returned by the `mysql_insert_id()` C API function.

- `lc_messages`

Command-Line Format	`--lc-messages=name`	
System Variable	**Name**	`lc_messages`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`en_US`

The locale to use for error messages. The default is `en_US`. The server converts the argument to a language name and combines it with the value of `lc_messages_dir` to produce the location for the error message file. See Section 10.2, “Setting the Error Message Language”.

- `lc_messages_dir`

Command-Line Format	`--lc-messages-dir=dir_name`	
System Variable	**Name**	`lc_messages_dir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The directory where error messages are located. The server uses the value together with the value of `lc_messages` to produce the location for the error message file. See Section 10.2, “Setting the Error Message Language”.

- `lc_time_names`

System Variable	**Name**	`lc_time_names`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string

This variable specifies the locale that controls the language used to display day and month names and abbreviations. This variable affects the output from the `DATE_FORMAT()`, `DAYNAME()` and `MONTHNAME()` functions. Locale names are POSIX-style values such as `'ja_JP'` or `'pt_BR'`. The default value is `'en_US'` regardless of your system's locale setting. For further information, see Section 10.7, “MySQL Server Locale Support”.

- `license`

System Variable	Name	license
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	GPL

The type of license the server has.

- `local_infile`

System Variable	Name	local_infile
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (<= 8.0.1)	Type	boolean
	Default	ON
Permitted Values (>= 8.0.2)	Type	boolean
	Default	OFF

This variable controls server-side `LOCAL` capability for `LOAD DATA` statements. Depending on the `local_infile` setting, the server refuses or permits local data loading by clients that have `LOCAL` enabled on the client side.

To explicitly cause the server to refuse or permit `LOAD DATA LOCAL` statements (regardless of how client programs and libraries are configured at build time or runtime), start `mysqld` with `local_infile` disabled or enabled, respectively. `local_infile` can also be set at runtime. For more information, see Section 6.1.6, “Security Issues with LOAD DATA LOCAL”.

- `lock_wait_timeout`

Command-Line Format	--lock-wait-timeout=#	
System Variable	Name	lock_wait_timeout
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	31536000
	Min Value	1

	Max Value	31536000

This variable specifies the timeout in seconds for attempts to acquire metadata locks. The permissible values range from 1 to 31536000 (1 year). The default is 31536000.

This timeout applies to all statements that use metadata locks. These include DML and DDL operations on tables, views, stored procedures, and stored functions, as well as `LOCK TABLES`, `FLUSH TABLES WITH READ LOCK`, and `HANDLER` statements.

This timeout does not apply to implicit accesses to system tables in the `mysql` database, such as grant tables modified by `GRANT` or `REVOKE` statements or table logging statements. The timeout does apply to system tables accessed directly, such as with `SELECT` or `UPDATE`.

The timeout value applies separately for each metadata lock attempt. A given statement can require more than one lock, so it is possible for the statement to block for longer than the `lock_wait_timeout` value before reporting a timeout error. When lock timeout occurs, `ER_LOCK_WAIT_TIMEOUT` is reported.

`lock_wait_timeout` does not apply to delayed inserts, which always execute with a timeout of 1 year. This is done to avoid unnecessary timeouts because a session that issues a delayed insert receives no notification of delayed insert timeouts.

- `locked_in_memory`

System Variable	**Name**	`locked_in_memory`
	Variable Scope	Global
	Dynamic Variable	No

Whether `mysqld` was locked in memory with `--memlock`.

- `log_bin_trust_function_creators`

Command-Line Format	`--log-bin-trust-function-creators`	
System Variable	**Name**	`log_bin_trust_function_creators`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`FALSE`

This variable applies when binary logging is enabled. It controls whether stored function creators can be trusted not to create stored functions that will cause unsafe events to be written to the binary log. If set to 0 (the default), users are not permitted to create or alter stored functions unless they have the `SUPER` privilege in addition to the `CREATE ROUTINE` or `ALTER ROUTINE` privilege. A setting of 0 also enforces the restriction that a function must be declared with the `DETERMINISTIC` characteristic, or with the `READS SQL DATA` or `NO SQL` characteristic. If the variable is set to 1, MySQL does not enforce these restrictions on stored function creation. This variable also applies to trigger creation. See Section 23.7, “Binary Logging of Stored Programs”.

- `log_builtin_as_identified_by_password`

Command-Line Format	`--log-builtin-as-identified-by-password[={OFF|ON}]`

System Variable	Name	log_builtin_as_identified_by_password
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

This variable affects binary logging of user-management statements. If enabled, binary logging for CREATE USER statements involving built-in authentication plugins rewrites the statements to include an IDENTIFIED BY PASSWORD clause, and SET PASSWORD statements are logged as SET PASSWORD statements, rather than being rewritten to ALTER USER statements.

- log_error

Command-Line Format	--log-error[=file_name]	
System Variable	Name	log_error
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The location of the error log, or stderr if the server is writing error message to the standard error output. See Section 5.4.2, “The Error Log”.

- log_error_verbosity

Command-Line Format	--log-error-verbosity=#	
System Variable	Name	log_error_verbosity
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	3
	Min Value	1
	Max Value	3

This variable controls verbosity of the server in writing error, warning, and note messages to the error log. The following table shows the permitted values. The default is 3.

Verbosity Value	Message Types Logged
1	Errors only
2	Errors and warnings
3	Errors, warnings, and notes

log_error_verbosity is preferred over, and should be used instead of, the older log_warnings system variable. See the description of log_warnings for information about how

that variable relates to `log_error_verbosity`. In particular, assigning a value to `log_warnings` assigns a value to `log_error_verbosity` and vice versa.

- `log_output`

Command-Line Format	`--log-output=name`	
System Variable	Name	`log_output`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	set
	Default	`FILE`
	Valid Values	`TABLE`
		`FILE`
		`NONE`

The destination for general query log and slow query log output. The value can be a comma-separated list of one or more of the words `TABLE` (log to tables), `FILE` (log to files), or `NONE` (do not log to tables or files). The default value is `FILE`. `NONE`, if present, takes precedence over any other specifiers. If the value is `NONE` log entries are not written even if the logs are enabled. If the logs are not enabled, no logging occurs even if the value of `log_output` is not `NONE`. For more information, see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”.

- `log_queries_not_using_indexes`

Command-Line Format	`--log-queries-not-using-indexes`	
System Variable	Name	`log_queries_not_using_indexes`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Whether queries that do not use indexes are logged to the slow query log. See Section 5.4.5, “The Slow Query Log”.

- `log_syslog`

Command-Line Format	`--log-syslog[={0|1}]`	
System Variable	Name	`log_syslog`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (Unix)	Type	boolean
	Default	`OFF`
Permitted Values (Windows)	Type	boolean
	Default	`ON`

Whether to write error log output to `syslog` (on Unix and Unix-like systems) or Event Log (on Windows). The default value is platform specific:

- On Unix and Unix-like systems, `syslog` output is disabled by default.
- On Windows, Event Log output is enabled by default, which is consistent with older MySQL versions.

Regardless of the default, `log_syslog` can be set explicitly to control output on any supported platform.

`syslog` output control is distinct from to sending error output to a file or (on Windows) to the console. Error output can be directed to the latter destination in addition to or instead of `syslog` as desired. See Section 5.4.2, “The Error Log”.

- `log_syslog_facility`

Command-Line Format	`--log-syslog-facility=value`	
System Variable	**Name**	`log_syslog_facility`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`daemon`

The facility for error log output written to `syslog` (what type of program is sending the message). This variable has no effect unless the `log_syslog` system variable is enabled. See Section 5.4.2, “The Error Log”.

The permitted values can vary per operating system; consult your system `syslog` documentation.

This variable does not exist on Windows.

- `log_syslog_include_pid`

Command-Line Format	`--log-syslog-include-pid[={0|1}]`	
System Variable	**Name**	`log_syslog_include_pid`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Whether to include the server process ID in each line of error log output written to `syslog`. This variable has no effect unless the `log_syslog` system variable is enabled. See Section 5.4.2, “The Error Log”.

This variable does not exist on Windows.

- `log_syslog_tag`

Command-Line Format	`--log-syslog-tag=value`	
System Variable	**Name**	`log_syslog_tag`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`empty string`

The tag to be added to the server identifier in error log output written to `syslog`. This variable has no effect unless the `log_syslog` system variable is enabled. See Section 5.4.2, “The Error Log”.

By default, the server identifier is `mysqld` with no tag. If a tag of *`tag_val`* is specified, it is appended to the server identifier with a leading hyphen, resulting in an identifier of `mysqld-`*`tag_val`*.

On Windows, to use a tag that does not already exist, the server must be run from an account with Administrator privileges, to permit creation of a registry entry for the tag. Elevated privileges are not required if the tag already exists.

- `log_timestamps`

Command-Line Format	`--log-timestamps=#`	
System Variable	**Name**	`log_timestamps`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`UTC`
	Valid Values	`UTC`
		`SYSTEM`

This variable controls the timestamp time zone of error log messages, and of general query log and slow query log messages written to files. It does not affect the time zone of general query log and slow query log messages written to tables (`mysql.general_log`, `mysql.slow_log`). Rows retrieved from those tables can be converted from the local system time zone to any desired time zone with `CONVERT_TZ()` or by setting the session `time_zone` system variable.

Permitted `log_timestamps` values are `UTC` (the default) and `SYSTEM` (local system time zone).

Timestamps are written using ISO 8601 / RFC 3339 format: `YYYY-MM-DDThh:mm:ss.uuuuuu` plus a tail value of `Z` signifying Zulu time (UTC) or `±hh:mm` (an offset from UTC).

- `log_throttle_queries_not_using_indexes`

System Variable	**Name**	`log_throttle_queries_not_using_indexes`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`

If `log_queries_not_using_indexes` is enabled, the `log_throttle_queries_not_using_indexes` variable limits the number of such queries per

minute that can be written to the slow query log. A value of 0 (the default) means “no limit”. For more information, see Section 5.4.5, “The Slow Query Log”.

- `log_slow_admin_statements`

System Variable	Name	log_slow_admin_statements
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Include slow administrative statements in the statements written to the slow query log. Administrative statements include `ALTER TABLE`, `ANALYZE TABLE`, `CHECK TABLE`, `CREATE INDEX`, `DROP INDEX`, `OPTIMIZE TABLE`, and `REPAIR TABLE`.

- `log_warnings`

Deprecated	5.7.2	
Command-Line Format	--log-warnings[=#]	
System Variable	Name	log_warnings
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	2
	Min Value	0
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	2
	Min Value	0
	Max Value	18446744073709551615

Whether to produce additional warning messages to the error log. This variable is enabled by default with a value of 2. To disable it, set it to 0. The server logs messages about statements that are unsafe for statement-based logging if the value is greater than 0. Aborted connections and access-denied errors for new connection attempts are logged if the value is greater than 1. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

Enabling this option by setting it greater than 0 is recommended, if you use replication, to get more information about what is happening, such as messages about network failures and reconnections. If the value is greater than 1, aborted connections are written to the error log, and access-denied errors for new connection attempts are written.

If a slave server is started with `log_warnings` enabled, the slave prints messages to the error log to provide information about its status, such as the binary log and relay log coordinates where it

starts its job, when it is switching to another relay log, when it reconnects after a disconnect, and so forth.

As of MySQL 5.7.2, information items previously governed by `log_warnings` are governed by `log_error_verbosity`, which is preferred over, and should be used instead of, the older `log_warnings` system variable. (The `log_warnings` system variable and `--log-warnings` command-line option are deprecated and will be removed in a future MySQL release.)

Assigning a value to `log_warnings` assigns a value to `log_error_verbosity` and vice versa. The variables are related as follows:

- Suppression of all `log_warnings` items, achieved with `log_warnings=0`, is achieved with `log_error_verbosity=1` (errors only).
- Items printed for `log_warnings=1` or higher count as warnings and are printed for `log_error_verbosity=2` or higher.
- Items printed for `log_warnings=2` count as notes and are printed for `log_error_verbosity=3`.

As of MySQL 5.7.2, the default log level is controlled by `log_error_verbosity`, which has a default of 3. In addition, the default for `log_warnings` changes from 1 to 2, which corresponds to `log_error_verbosity=3`. To achieve a logging level similar to the previous default, set `log_error_verbosity=2`.

In MySQL 5.7.2 and higher, use of `log_warnings` is still permitted but maps onto use of `log_error_verbosity` as follows:

- Setting `log_warnings=0` is equivalent to `log_error_verbosity=1` (errors only).
- Setting `log_warnings=1` is equivalent to `log_error_verbosity=2` (errors, warnings).
- Setting `log_warnings=2` (or higher) is equivalent to `log_error_verbosity=3` (errors, warnings, notes), and the server sets `log_warnings` to 2 if a larger value is specified.

- `long_query_time`

Command-Line Format	`--long-query-time=#`	
System Variable	**Name**	`long_query_time`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	numeric
	Default	`10`
	Min Value	`0`

If a query takes longer than this many seconds, the server increments the `Slow_queries` status variable. If the slow query log is enabled, the query is logged to the slow query log file. This value is measured in real time, not CPU time, so a query that is under the threshold on a lightly loaded system might be above the threshold on a heavily loaded one. The minimum and default values of `long_query_time` are 0 and 10, respectively. The value can be specified to a resolution of microseconds. For logging to a file, times are written including the microseconds part. For logging to tables, only integer times are written; the microseconds part is ignored. See Section 5.4.5, “The Slow Query Log”.

- `low_priority_updates`

Command-Line Format	--low-priority-updates	
System Variable	Name	low_priority_updates
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	FALSE

If set to 1, all INSERT, UPDATE, DELETE, and LOCK TABLE WRITE statements wait until there is no pending SELECT or LOCK TABLE READ on the affected table. This affects only storage engines that use only table-level locking (such as MyISAM, MEMORY, and MERGE).

- lower_case_file_system

System Variable	Name	lower_case_file_system
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean

This variable describes the case sensitivity of file names on the file system where the data directory is located. OFF means file names are case sensitive, ON means they are not case sensitive. This variable is read only because it reflects a file system attribute and setting it would have no effect on the file system.

- lower_case_table_names

Command-Line Format	--lower-case-table-names[=#]	
System Variable	Name	lower_case_table_names
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	0
	Min Value	0
	Max Value	2

If set to 0, table names are stored as specified and comparisons are case sensitive. If set to 1, table names are stored in lowercase on disk and comparisons are not case sensitive. If set to 2, table names are stored as given but compared in lowercase. This option also applies to database names and table aliases. For additional information, see Section 9.2.2, “Identifier Case Sensitivity”.

On Windows the default value is 1. On OS X, the default value is 2.

You should *not* set lower_case_table_names to 0 if you are running MySQL on a system where the data directory resides on a case-insensitive file system (such as on Windows or OS X). It is an unsupported combination that could result in a hang condition when running an INSERT INTO ... .

`SELECT ... FROM tbl_name` operation with the wrong *`tbl_name`* letter case. With `MyISAM`, accessing table names using different letter cases could cause index corruption.

As of MySQL 5.7.9, an error message is printed and the server exits if you attempt to start the server with `--lower_case_table_names=0` on a case-insensitive file system.

If you are using `InnoDB` tables, you should set this variable to 1 on all platforms to force names to be converted to lowercase.

The setting of this variable in MySQL 8.0 affects the behavior of replication filtering options with regard to case sensitivity. (Bug #51639) See Section 18.2.5, "How Servers Evaluate Replication Filtering Rules", for more information.

- `max_allowed_packet`

Command-Line Format	`--max-allowed-packet=#`	
System Variable	**Name**	`max_allowed_packet`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`4194304`
	Min Value	`1024`
	Max Value	`1073741824`

The maximum size of one packet or any generated/intermediate string, or any parameter sent by the `mysql_stmt_send_long_data()` C API function. The default is 4MB.

The packet message buffer is initialized to `net_buffer_length` bytes, but can grow up to `max_allowed_packet` bytes when needed. This value by default is small, to catch large (possibly incorrect) packets.

You must increase this value if you are using large `BLOB` columns or long strings. It should be as big as the largest `BLOB` you want to use. The protocol limit for `max_allowed_packet` is 1GB. The value should be a multiple of 1024; nonmultiples are rounded down to the nearest multiple.

When you change the message buffer size by changing the value of the `max_allowed_packet` variable, you should also change the buffer size on the client side if your client program permits it. The default `max_allowed_packet` value built in to the client library is 1GB, but individual client programs might override this. For example, `mysql` and `mysqldump` have defaults of 16MB and 24MB, respectively. They also enable you to change the client-side value by setting `max_allowed_packet` on the command line or in an option file.

The session value of this variable is read only. The client can receive up to as many bytes as the session value. However, the server will not send to the client more bytes than the current global `max_allowed_packet` value. (The global value could be less than the session value if the global value is changed after the client connects.)

- `max_connect_errors`

Command-Line Format	`--max-connect-errors=#`	
System Variable	**Name**	`max_connect_errors`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`100`
	Min Value	`1`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`100`
	Min Value	`1`
	Max Value	`18446744073709551615`

If more than this many successive connection requests from a host are interrupted without a successful connection, the server blocks that host from further connections. You can unblock blocked hosts by flushing the host cache. To do so, issue a `FLUSH HOSTS` statement or execute a `mysqladmin flush-hosts` command. If a connection is established successfully within fewer than `max_connect_errors` attempts after a previous connection was interrupted, the error count for the host is cleared to zero. However, once a host is blocked, flushing the host cache is the only way to unblock it. The default is 100.

- `max_connections`

Command-Line Format	`--max-connections=#`	
System Variable	**Name**	`max_connections`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`151`
	Min Value	`1`
	Max Value	`100000`

The maximum permitted number of simultaneous client connections. By default, this is 151. See Section B.5.2.6, “Too many connections”, for more information.

Increasing this value increases the number of file descriptors that `mysqld` requires. If the required number of descriptors are not available, the server reduces the value of `max_connections`. See Section 8.4.3.1, “How MySQL Opens and Closes Tables”, for comments on file descriptor limits.

`mysqld` actually permits `max_connections+1` clients to connect. The extra connection is reserved for use by accounts that have the `SUPER` privilege. By granting the `SUPER` privilege to administrators and not to normal users (who should not need it), an administrator can connect to the server and use `SHOW PROCESSLIST` to diagnose problems even if the maximum number of unprivileged clients are connected. See Section 13.7.5.29, “SHOW PROCESSLIST Syntax”.

Connections refused because the `max_connections` limit is reached increment the `Connection_errors_max_connections` status variable.

- `max_delayed_threads`

Deprecated	5.6.7	
Command-Line Format	`--max-delayed-threads=#`	
System Variable	**Name**	`max_delayed_threads`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`20`
	Min Value	`0`
	Max Value	`16384`

This system variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `max_digest_length`

Command-Line Format	`--max-digest-length=#`	
System Variable	**Name**	`max_digest_length`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`1024`
	Min Value	`0`
	Max Value	`1048576`

The maximum number of bytes available for computing normalized statement digests. Once that amount of space is used during digest computation, truncation occurs: no further tokens from a parsed statement are collected or figure into its digest value. Statements that differ only after that many bytes of parsed tokens produce the same normalized statement digest and are considered identical if compared or if aggregated for digest statistics.

Decreasing the `max_digest_length` value reduces memory use but causes the digest value of more statements to become indistinguishable if they differ only at the end. Increasing the value permits longer statements to be distinguished but increases memory use, particularly for workloads that involve large numbers of simultaneous sessions (the server allocates `max_digest_length` bytes per session).

The parser uses this system variable as a limit on the maximum length of normalized statement digests that it computes. The Performance Schema, if it tracks statement digests, makes a copy of the digest value, using the `performance_schema_max_digest_length`.

system variable as a limit on the maximum length of digests that it stores. Consequently, if `performance_schema_max_digest_length` is less than `max_digest_length`, digest values stored in the Performance Schema are truncated relative to the original digest values.

For more information about statement digesting, see Section 25.9, "Performance Schema Statement Digests".

- `max_error_count`

Command-Line Format	`--max-error-count=#`	
System Variable	Name	`max_error_count`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (<= 8.0.2)	Type	integer
	Default	`64`
	Min Value	`0`
	Max Value	`65535`
Permitted Values (>= 8.0.3)	Type	integer
	Default	`1024`
	Min Value	`0`
	Max Value	`65535`

The maximum number of error, warning, and note messages to be stored for display by the `SHOW ERRORS` and `SHOW WARNINGS` statements. This is the same as the number of condition areas in the diagnostics area, and thus the number of conditions that can be inspected by `GET DIAGNOSTICS`.

- `max_execution_time`

Command-Line Format	`--max-execution-time=#`	
System Variable	Name	`max_execution_time`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`

The execution timeout for `SELECT` statements, in milliseconds. If the value is 0, timeouts are not enabled.

`max_execution_time` applies as follows:

- The global `max_execution_time` value provides the default for the session value for new connections. The session value applies to `SELECT` executions executed within the session that include no `MAX_EXECUTION_TIME(N)` optimizer hint or for which *`N`* is 0.

- `max_execution_time` applies to read-only `SELECT` statements. Statements that are not read only are those that invoke a stored function that modifies data as a side effect.
- `max_execution_time` is ignored for `SELECT` statements in stored programs.

- `max_heap_table_size`

Command-Line Format	`--max-heap-table-size=#`	
System Variable	**Name**	`max_heap_table_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`16777216`
	Min Value	`16384`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`16777216`
	Min Value	`16384`
	Max Value	`18446744073709551615`

This variable sets the maximum size to which user-created `MEMORY` tables are permitted to grow. The value of the variable is used to calculate `MEMORY` table `MAX_ROWS` values. Setting this variable has no effect on any existing `MEMORY` table, unless the table is re-created with a statement such as `CREATE TABLE` or altered with `ALTER TABLE` or `TRUNCATE TABLE`. A server restart also sets the maximum size of existing `MEMORY` tables to the global `max_heap_table_size` value.

This variable is also used in conjunction with `tmp_table_size` to limit the size of internal in-memory tables. See Section 8.4.4, “Internal Temporary Table Use in MySQL”.

`max_heap_table_size` is not replicated. See Section 18.4.1.23, “Replication and MEMORY Tables”, and Section 18.4.1.38, “Replication and Variables”, for more information.

- `max_insert_delayed_threads`

Deprecated	5.6.7	
System Variable	**Name**	`max_insert_delayed_threads`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer

This variable is a synonym for `max_delayed_threads`.

This system variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `max_join_size`

Command-Line Format	`--max-join-size=#`	
System Variable	Name	`max_join_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`18446744073709551615`
	Min Value	`1`
	Max Value	`18446744073709551615`

Do not permit statements that probably need to examine more than `max_join_size` rows (for single-table statements) or row combinations (for multiple-table statements) or that are likely to do more than `max_join_size` disk seeks. By setting this value, you can catch statements where keys are not used properly and that would probably take a long time. Set it if your users tend to perform joins that lack a `WHERE` clause, that take a long time, or that return millions of rows.

Setting this variable to a value other than `DEFAULT` resets the value of `sql_big_selects` to `0`. If you set the `sql_big_selects` value again, the `max_join_size` variable is ignored.

If a query result is in the query cache, no result size check is performed, because the result has previously been computed and it does not burden the server to send it to the client.

- `max_length_for_sort_data`

Command-Line Format	`--max-length-for-sort-data=#`	
System Variable	Name	`max_length_for_sort_data`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1024`
	Min Value	`4`
	Max Value	`8388608`

The cutoff on the size of index values that determines which `filesort` algorithm to use. See Section 8.2.1.13, “ORDER BY Optimization”.

- `max_points_in_geometry`

Command-Line Format	`--max-points-in-geometry=integer`	
System Variable	Name	`max_points_in_geometry`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`65536`
	Min Value	`3`
	Max Value	`1048576`

The maximum value of the *`points_per_circle`* argument to the `ST_Buffer_Strategy()` function.

- `max_prepared_stmt_count`

Command-Line Format	`--max-prepared-stmt-count=#`	
System Variable	**Name**	`max_prepared_stmt_count`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`16382`
	Min Value	`0`
	Max Value	`1048576`

This variable limits the total number of prepared statements in the server. It can be used in environments where there is the potential for denial-of-service attacks based on running the server out of memory by preparing huge numbers of statements. If the value is set lower than the current number of prepared statements, existing statements are not affected and can be used, but no new statements can be prepared until the current number drops below the limit. The default value is 16,382. The permissible range of values is from 0 to 1 million. Setting the value to 0 disables prepared statements.

- `max_relay_log_size`

Command-Line Format	`--max-relay-log-size=#`	
System Variable	**Name**	`max_relay_log_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1073741824`

If a write by a replication slave to its relay log causes the current log file size to exceed the value of this variable, the slave rotates the relay logs (closes the current file and opens the next one). If `max_relay_log_size` is 0, the server uses `max_binlog_size` for both the binary log and the relay log. If `max_relay_log_size` is greater than 0, it constrains the size of the relay log, which enables you to have different sizes for the two logs. You must set `max_relay_log_size` to between 4096 bytes and 1GB (inclusive), or to 0. The default value is 0. See Section 18.2.2, "Replication Implementation Details".

- `max_seeks_for_key`

Command-Line Format	`--max-seeks-for-key=#`	
System Variable	**Name**	`max_seeks_for_key`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`4294967295`
	Min Value	`1`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`18446744073709551615`
	Min Value	`1`
	Max Value	`18446744073709551615`

Limit the assumed maximum number of seeks when looking up rows based on a key. The MySQL optimizer assumes that no more than this number of key seeks are required when searching for matching rows in a table by scanning an index, regardless of the actual cardinality of the index (see Section 13.7.5.22, "SHOW INDEX Syntax"). By setting this to a low value (say, 100), you can force MySQL to prefer indexes instead of table scans.

- `max_sort_length`

Command-Line Format	`--max-sort-length=#`	
System Variable	**Name**	`max_sort_length`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1024`
	Min Value	`4`
	Max Value	`8388608`

The number of bytes to use when sorting data values. The server uses only the first `max_sort_length` bytes of each value and ignores the rest. Consequently, values that differ only after the first `max_sort_length` bytes compare as equal for `GROUP BY`, `ORDER BY`, and `DISTINCT` operations.

Increasing the value of `max_sort_length` may require increasing the value of `sort_buffer_size` as well. For details, see Section 8.2.1.13, “ORDER BY Optimization”

- `max_sp_recursion_depth`

Command-Line Format	`--max-sp-recursion-depth[=#]`	
System Variable	Name	`max_sp_recursion_depth`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Max Value	`255`

The number of times that any given stored procedure may be called recursively. The default value for this option is 0, which completely disables recursion in stored procedures. The maximum value is 255.

Stored procedure recursion increases the demand on thread stack space. If you increase the value of `max_sp_recursion_depth`, it may be necessary to increase thread stack size by increasing the value of `thread_stack` at server startup.

- `max_tmp_tables`

This variable is unused. It is deprecated and will be removed in a future MySQL release.

- `max_user_connections`

Command-Line Format	`--max-user-connections=#`	
System Variable	Name	`max_user_connections`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

The maximum number of simultaneous connections permitted to any given MySQL user account. A value of 0 (the default) means “no limit.”

This variable has a global value that can be set at server startup or runtime. It also has a read-only session value that indicates the effective simultaneous-connection limit that applies to the account associated with the current session. The session value is initialized as follows:

- If the user account has a nonzero `MAX_USER_CONNECTIONS` resource limit, the session `max_user_connections` value is set to that limit.
- Otherwise, the session `max_user_connections` value is set to the global value.

Account resource limits are specified using the `CREATE USER` or `ALTER USER` statement. See Section 6.3.6, “Setting Account Resource Limits”.

- `max_write_lock_count`

Command-Line Format	`--max-write-lock-count=#`	
System Variable	**Name**	`max_write_lock_count`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`4294967295`
	Min Value	`1`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`18446744073709551615`
	Min Value	`1`
	Max Value	`18446744073709551615`

After this many write locks, permit some pending read lock requests to be processed in between.

- `mecab_rc_file`

Command-Line Format	`--mecab-rc-file`	
System Variable	**Name**	`mecab_rc_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The `mecab_rc_file` option is used when setting up the MeCab full-text parser.

The `mecab_rc_file` option defines the path to the `mecabrc` configuration file, which is the configuration file for MeCab. The option is read-only and can only be set at startup. The `mecabrc` configuration file is required to initialize MeCab.

For information about the MeCab full-text parser, see Section 12.9.9, “MeCab Full-Text Parser Plugin”.

For information about options that can be specified in the MeCab `mecabrc` configuration file, refer to the MeCab Documentation on the Google Developers site.

- `metadata_locks_cache_size`

Deprecated	5.7.4	
System Variable	Name	metadata_locks_cache_size
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	1024
	Min Value	1
	Max Value	1048576

The size of the metadata locks cache. The server uses this cache to avoid creation and destruction of synchronization objects. This is particularly helpful on systems where such operations are expensive, such as Windows XP.

Metadata locking implementation changes have made this variable unnecessary, so it is deprecated and will be removed in a future MySQL release.

- `metadata_locks_hash_instances`

Deprecated	5.7.4	
System Variable	Name	metadata_locks_hash_instances
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	8
	Min Value	1
	Max Value	1024

The set of metadata locks can be partitioned into separate hashes to permit connections accessing different objects to use different locking hashes and reduce contention. The `metadata_locks_hash_instances` system variable specifies the number of hashes (default 8).

Metadata locking implementation changes have made this variable unnecessary, so it is deprecated and will be removed in a future MySQL release.

- `min_examined_row_limit`

Command-Line Format	--min-examined-row-limit=#	
System Variable	Name	min_examined_row_limit
	Variable Scope	Global, Session
	Dynamic Variable	Yes

Permitted Values (32-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709551615`

Queries that examine fewer than this number of rows are not logged to the slow query log.

- `multi_range_count`

Deprecated	5.6.7	
Command-Line Format	`--multi-range-count=#`	
System Variable	**Name**	`multi_range_count`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`256`
	Min Value	`1`
	Max Value	`4294967295`

This variable has no effect. It is deprecated and will be removed in a future MySQL release.

- `myisam_data_pointer_size`

Command-Line Format	`--myisam-data-pointer-size=#`	
System Variable	**Name**	`myisam_data_pointer_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`6`
	Min Value	`2`
	Max Value	`7`

The default pointer size in bytes, to be used by `CREATE TABLE` for `MyISAM` tables when no `MAX_ROWS` option is specified. This variable cannot be less than 2 or larger than 7. The default value is 6. See Section B.5.2.11, “The table is full”.

- `myisam_max_sort_file_size`

Command-Line Format	`--myisam-max-sort-file-size=#`	
System Variable	**Name**	`myisam_max_sort_file_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`2147483648`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`9223372036854775807`

The maximum size of the temporary file that MySQL is permitted to use while re-creating a `MyISAM` index (during `REPAIR TABLE`, `ALTER TABLE`, or `LOAD DATA INFILE`). If the file size would be larger than this value, the index is created using the key cache instead, which is slower. The value is given in bytes.

If `MyISAM` index files exceed this size and disk space is available, increasing the value may help performance. The space must be available in the file system containing the directory where the original index file is located.

- `myisam_mmap_size`

Command-Line Format	`--myisam-mmap-size=#`	
System Variable	**Name**	`myisam_mmap_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`4294967295`
	Min Value	`7`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`18446744073709551615`
	Min Value	`7`
	Max Value	`18446744073709551615`

The maximum amount of memory to use for memory mapping compressed `MyISAM` files. If many compressed `MyISAM` tables are used, the value can be decreased to reduce the likelihood of memory-swapping problems.

- `myisam_recover_options`

System Variable	Name	myisam_recover_options
	Variable Scope	Global
	Dynamic Variable	No

The value of the `--myisam-recover-options` option. See Section 5.1.4, “Server Command Options”.

- `myisam_repair_threads`

Command-Line Format	--myisam-repair-threads=#	
System Variable	Name	myisam_repair_threads
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	1
	Min Value	1
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	1
	Min Value	1
	Max Value	18446744073709551615

If this value is greater than 1, `MyISAM` table indexes are created in parallel (each index in its own thread) during the `Repair by sorting` process. The default value is 1.

Note

Multi-threaded repair is still *beta-quality* code.

- `myisam_sort_buffer_size`

Command-Line Format	--myisam-sort-buffer-size=#	
System Variable	Name	myisam_sort_buffer_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (Windows, 32-bit platforms)	Type	integer
	Default	8388608
	Min Value	4096

	Max Value	4294967295
Permitted Values (Windows, 64-bit platforms)	**Type**	integer
	Default	8388608
	Min Value	4096
	Max Value	18446744073709551615
Permitted Values (Other, 32-bit platforms)	**Type**	integer
	Default	8388608
	Min Value	4096
	Max Value	4294967295
Permitted Values (Other, 64-bit platforms)	**Type**	integer
	Default	8388608
	Min Value	4096
	Max Value	18446744073709551615

The size of the buffer that is allocated when sorting `MyISAM` indexes during a `REPAIR TABLE` or when creating indexes with `CREATE INDEX` or `ALTER TABLE`.

The maximum permissible setting for `myisam_sort_buffer_size` is 4GB−1. Larger values are permitted for 64-bit platforms.

- `myisam_stats_method`

Command-Line Format	`--myisam-stats-method=name`	
System Variable	**Name**	`myisam_stats_method`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`nulls_unequal`
	Valid Values	`nulls_equal`
		`nulls_unequal`
		`nulls_ignored`

How the server treats `NULL` values when collecting statistics about the distribution of index values for `MyISAM` tables. This variable has three possible values, `nulls_equal`, `nulls_unequal`, and `nulls_ignored`. For `nulls_equal`, all `NULL` index values are considered equal and form a single value group that has a size equal to the number of `NULL` values. For `nulls_unequal`, `NULL` values are considered unequal, and each `NULL` forms a distinct value group of size 1. For `nulls_ignored`, `NULL` values are ignored.

The method that is used for generating table statistics influences how the optimizer chooses indexes for query execution, as described in Section 8.3.7, “InnoDB and MyISAM Index Statistics Collection”.

- `myisam_use_mmap`

Command-Line Format	`--myisam-use-mmap`	
System Variable	Name	`myisam_use_mmap`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Use memory mapping for reading and writing `MyISAM` tables.

- `mysql_native_password_proxy_users`

Command-Line Format	`--mysql-native-password-proxy-users=[={OFF|ON}]`	
System Variable	Name	`mysql_native_password_proxy_users`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

This variable controls whether the `mysql_native_password` built-in authentication plugin supports proxy users. It has no effect unless the `check_proxy_users` system variable is enabled. For information about user proxying, see Section 6.3.11, “Proxy Users”.

- `named_pipe`

System Variable	Name	`named_pipe`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Windows	
Permitted Values (Windows)	Type	boolean
	Default	`OFF`

(Windows only.) Indicates whether the server supports connections over named pipes.

- `net_buffer_length`

Command-Line Format	`--net-buffer-length=#`	
System Variable	Name	`net_buffer_length`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`16384`

	Min Value	`1024`
	Max Value	`1048576`

Each client thread is associated with a connection buffer and result buffer. Both begin with a size given by `net_buffer_length` but are dynamically enlarged up to `max_allowed_packet` bytes as needed. The result buffer shrinks to `net_buffer_length` after each SQL statement.

This variable should not normally be changed, but if you have very little memory, you can set it to the expected length of statements sent by clients. If statements exceed this length, the connection buffer is automatically enlarged. The maximum value to which `net_buffer_length` can be set is 1MB.

The session value of this variable is read only.

- `net_read_timeout`

Command-Line Format	`--net-read-timeout=#`	
System Variable	**Name**	`net_read_timeout`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`30`
	Min Value	`1`

The number of seconds to wait for more data from a connection before aborting the read. When the server is reading from the client, `net_read_timeout` is the timeout value controlling when to abort. When the server is writing to the client, `net_write_timeout` is the timeout value controlling when to abort. See also `slave_net_timeout`.

- `net_retry_count`

Command-Line Format	`--net-retry-count=#`	
System Variable	**Name**	`net_retry_count`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`10`
	Min Value	`1`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`10`
	Min Value	`1`

	Max Value	18446744073709551615

If a read or write on a communication port is interrupted, retry this many times before giving up. This value should be set quite high on FreeBSD because internal interrupts are sent to all threads.

- net_write_timeout

Command-Line Format	--net-write-timeout=#	
System Variable	**Name**	net_write_timeout
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	60
	Min Value	1

The number of seconds to wait for a block to be written to a connection before aborting the write. See also net_read_timeout.

- new

Command-Line Format	--new	
System Variable	**Name**	new
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Disabled by	skip-new	
Permitted Values	**Type**	boolean
	Default	FALSE

This variable was used in MySQL 4.0 to turn on some 4.1 behaviors, and is retained for backward compatibility. Its value is always OFF.

- ngram_token_size

Command-Line Format	--ngram-token-size	
System Variable	**Name**	ngram_token_size
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	2
	Min Value	1
	Max Value	10

Defines the n-gram token size for the n-gram full-text parser. The `ngram_token_size` option is read-only and can only be modified at startup. The default value is 2 (bigram). The maximum value is 10.

For more information about how to configure this variable, see Section 12.9.8, “ngram Full-Text Parser”.

- `offline_mode`

Command-Line Format	`--offline-mode=val`	
System Variable	Name	`offline_mode`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Whether the server is in “offline mode”, which has these characteristics:

- Connected client users who do not have the `CONNECTION_ADMIN` or `SUPER` privilege are disconnected on the next request, with an appropriate error. Disconnection includes terminating running statements and releasing locks. Such clients also cannot initiate new connections, and receive an appropriate error.
- Connected client users who have the `CONNECTION_ADMIN` or `SUPER` privilege are not disconnected, and can initiate new connections to manage the server.
- Replication slave threads are permitted to keep applying data to the server.

Only users who have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege can control offline mode. To put a server in offline mode, change the value of the `offline_mode` system variable from `OFF` to `ON`. To resume normal operations, change `offline_mode` from `ON` to `OFF`. In offline mode, clients that are refused access receive an `ER_SERVER_OFFLINE_MODE` error.

- `old`

Command-Line Format	`--old`	
System Variable	Name	`old`
	Variable Scope	Global
	Dynamic Variable	No

`old` is a compatibility variable. It is disabled by default, but can be enabled at startup to revert the server to behaviors present in older versions.

When `old` is enabled, it changes the default scope of index hints to that used prior to MySQL 5.1.17. That is, index hints with no `FOR` clause apply only to how indexes are used for row retrieval and not to resolution of `ORDER BY` or `GROUP BY` clauses. (See Section 8.9.4, “Index Hints”.) Take care about enabling this in a replication setup. With statement-based binary logging, having different modes for the master and slaves might lead to replication errors.

- `old_alter_table`

Command-Line Format	`--old-alter-table`

System Variable	Name	old_alter_table
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

When this variable is enabled, the server does not use the optimized method of processing an ALTER TABLE operation. It reverts to using a temporary table, copying over the data, and then renaming the temporary table to the original, as used by MySQL 5.0 and earlier. For more information on the operation of ALTER TABLE, see Section 13.1.7, “ALTER TABLE Syntax”.

As of MySQL 8.0.2, ALTER TABLE ... DROP PARTITION with old_alter_table=ON rebuilds the partitioned table and attempts to move data from the dropped partition to another partition with a compatible PARTITION ... VALUES definition. Data that cannot be moved to another partition is deleted. In earlier releases, ALTER TABLE ... DROP PARTITION with old_alter_table=ON deletes data stored in the partition and drops the partition.

- old_passwords

Deprecated	5.7.6	
System Variable	Name	old_passwords
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	0
	Valid Values	0
		2

Note

This system variable is deprecated and will be removed in a future MySQL release.

This variable controls the password hashing method used by the PASSWORD() function. It also influences password hashing performed by CREATE USER and GRANT statements that specify a password using an IDENTIFIED BY clause.

The following table shows the permitted values of old_passwords, the password hashing method for each value, and which authentication plugins use passwords hashed with each method.

Value	Password Hashing Method	Associated Authentication Plugin
0	MySQL 4.1 native hashing	mysql_native_password
2	SHA-256 hashing	sha256_password

If you set old_passwords=2, follow the instructions for using the sha256_password plugin at Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

The server sets the global old_passwords value during startup to be consistent with the password hashing method required by the default authentication plugin. The default plugin is

`mysql_native_password` unless the `default_authentication_plugin` system variable is set otherwise.

When a client successfully connects to the server, the server sets the session `old_passwords` value appropriately for the account authentication method. For example, if the account uses the `sha256_password` authentication plugin, the server sets `old_passwords=2`.

For additional information about authentication plugins, see Section 6.3.10, “Pluggable Authentication”.

- `open_files_limit`

Command-Line Format	`--open-files-limit=#`	
System Variable	Name	`open_files_limit`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`5000, with possible adjustment`
	Min Value	`0`
	Max Value	`platform dependent`

The number of files that the operating system permits `mysqld` to open. The value of this variable at runtime is the real value permitted by the system and might be different from the value you specify at server startup. The value is 0 on systems where MySQL cannot change the number of open files.

The effective `open_files_limit` value is based on the value specified at system startup (if any) and the values of `max_connections` and `table_open_cache`, using these formulas:

```
1) 10 + max_connections + (table_open_cache * 2)
2) max_connections * 5
3) open_files_limit value specified at startup, 5000 if none
```

The server attempts to obtain the number of file descriptors using the maximum of those three values. If that many descriptors cannot be obtained, the server attempts to obtain as many as the system will permit.

- `optimizer_prune_level`

Command-Line Format	`--optimizer-prune-level[=#]`	
System Variable	Name	`optimizer_prune_level`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`1`

Controls the heuristics applied during query optimization to prune less-promising partial plans from the optimizer search space. A value of 0 disables heuristics so that the optimizer performs an

exhaustive search. A value of 1 causes the optimizer to prune plans based on the number of rows retrieved by intermediate plans.

- `optimizer_search_depth`

Command-Line Format	`--optimizer-search-depth[=#]`	
System Variable	Name	`optimizer_search_depth`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`62`
	Min Value	`0`
	Max Value	`62`

The maximum depth of search performed by the query optimizer. Values larger than the number of relations in a query result in better query plans, but take longer to generate an execution plan for a query. Values smaller than the number of relations in a query return an execution plan quicker, but the resulting plan may be far from being optimal. If set to 0, the system automatically picks a reasonable value.

- `optimizer_switch`

Command-Line Format	`--optimizer-switch=value`	
System Variable	Name	`optimizer_switch`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	set
	Valid Values	`batched_key_access={on|off}`
		`block_nested_loop={on|off}`
		`condition_fanout_filter={on|off}`
		`derived_merge={on|off}`
		`duplicateweedout={on|off}`
		`engine_condition_pushdown={on|off}`
		`firstmatch={on|off}`
		`index_condition_pushdown={on|off}`
		`index_merge={on|off}`
		`index_merge_intersection={on|off}`
		`index_merge_sort_union={on|off}`
		`index_merge_union={on|off}`
		`loosescan={on|off}`
		`materialization={on|off}`
		`mrr={on|off}`

		mrr_cost_based={on|off}
		semijoin={on|off}
		subquery_materialization_cost_based={on|off}
		use_index_extensions={on|off}

The optimizer_switch system variable enables control over optimizer behavior. The value of this variable is a set of flags, each of which has a value of on or off to indicate whether the corresponding optimizer behavior is enabled or disabled. This variable has global and session values and can be changed at runtime. The global default can be set at server startup.

To see the current set of optimizer flags, select the variable value:

```
mysql> SELECT @@optimizer_switch\G
*************************** 1. row ***************************
@@optimizer_switch: index_merge=on,index_merge_union=on,
                    index_merge_sort_union=on,
                    index_merge_intersection=on,
                    engine_condition_pushdown=on,
                    index_condition_pushdown=on,
                    mrr=on,mrr_cost_based=on,
                    block_nested_loop=on,batched_key_access=off,
                    materialization=on,semijoin=on,loosescan=on,
                    firstmatch=on,duplicateweedout=on,
                    subquery_materialization_cost_based=on,
                    use_index_extensions=on,
                    condition_fanout_filter=on,derived_merge=on
```

For more information about the syntax of this variable and the optimizer behaviors that it controls, see Section 8.9.3, “Switchable Optimizations”.

- optimizer_trace

System Variable	Name	optimizer_trace
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string

This variable controls optimizer tracing. For details, see MySQL Internals: Tracing the Optimizer.

- optimizer_trace_features

System Variable	Name	optimizer_trace_features
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string

This variable enables or disables selected optimizer tracing features. For details, see MySQL Internals: Tracing the Optimizer.

- optimizer_trace_limit

System Variable	Name	optimizer_trace_limit

	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1`

The maximum number of optimizer traces to display. For details, see MySQL Internals: Tracing the Optimizer.

- `optimizer_trace_max_mem_size`

System Variable	**Name**	`optimizer_trace_max_mem_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`16384`

The maximum cumulative size of stored optimizer traces. For details, see MySQL Internals: Tracing the Optimizer.

- `optimizer_trace_offset`

System Variable	**Name**	`optimizer_trace_offset`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`-1`

The offset of optimizer traces to display. For details, see MySQL Internals: Tracing the Optimizer.

- `performance_schema_xxx`

Performance Schema system variables are listed in Section 25.14, “Performance Schema System Variables”. These variables may be used to configure Performance Schema operation.

- `parser_max_mem_size`

Command-Line Format	`--parser-max-mem-size=N`	
System Variable	**Name**	`parser_max_mem_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`4294967295`
	Min Value	`10000000`

	Max Value	4294967295
Permitted Values (64-bit platforms)	**Type**	integer
	Default	18446744073709551615
	Min Value	10000000
	Max Value	18446744073709551615

The maximum amount of memory available to the parser. The default value places no limit on memory available. The value can be reduced to protect against out-of-memory situations caused by parsing long or complex SQL statements.

- `persisted_globals_load`

Command-Line Format	`--persisted-globals-load[=ON|OFF]`	
System Variable	**Name**	`persisted_globals_load`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`ON`

Whether to load persisted configuration settings from the `mysqld-auto.cnf` file in the data directory. The server normally processes this file at startup after all other option files (see Section 4.2.6, “Using Option Files”). Disabling `persisted_globals_load` causes the server startup sequence to skip `mysqld-auto.cnf`.

To modify the contents of `mysqld-auto.cnf`, use the `SET PERSIST` statement. See Section 13.7.4.1, “SET Syntax for Variable Assignment”.

- `pid_file`

Command-Line Format	`--pid-file=file_name`	
System Variable	**Name**	`pid_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The path name of the process ID (PID) file. This variable can be set with the `--pid-file` option.

- `plugin_dir`

Command-Line Format	`--plugin-dir=dir_name`	
System Variable	**Name**	`plugin_dir`
	Variable Scope	Global
	Dynamic Variable	No

Permitted Values	Type	directory name
	Default	BASEDIR/lib/plugin

The path name of the plugin directory.

If the plugin directory is writable by the server, it may be possible for a user to write executable code to a file in the directory using SELECT ... INTO DUMPFILE. This can be prevented by making plugin_dir read only to the server or by setting --secure-file-priv to a directory where SELECT writes can be made safely.

- port

Command-Line Format	--port=#	
System Variable	Name	port
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	3306
	Min Value	0
	Max Value	65535

The number of the port on which the server listens for TCP/IP connections. This variable can be set with the --port option.

- preload_buffer_size

Command-Line Format	--preload-buffer-size=#	
System Variable	Name	preload_buffer_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	32768
	Min Value	1024
	Max Value	1073741824

The size of the buffer that is allocated when preloading indexes.

- profiling

If set to 0 or OFF (the default), statement profiling is disabled. If set to 1 or ON, statement profiling is enabled and the SHOW PROFILE and SHOW PROFILES statements provide access to profiling information. See Section 13.7.5.31, “SHOW PROFILES Syntax”.

This variable is deprecated and will be removed in a future MySQL release.

- profiling_history_size

The number of statements for which to maintain profiling information if `profiling` is enabled. The default value is 15. The maximum value is 100. Setting the value to 0 effectively disables profiling. See Section 13.7.5.31, “SHOW PROFILES Syntax”.

This variable is deprecated and will be removed in a future MySQL release.

- `protocol_version`

System Variable	Name	`protocol_version`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer

The version of the client/server protocol used by the MySQL server.

- `proxy_user`

System Variable	Name	`proxy_user`
	Variable Scope	Session
	Dynamic Variable	No
Permitted Values	Type	string

If the current client is a proxy for another user, this variable is the proxy user account name. Otherwise, this variable is `NULL`. See Section 6.3.11, “Proxy Users”.

- `pseudo_slave_mode`

System Variable	Name	`pseudo_slave_mode`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	Type	integer

This variable is for internal server use.

- `pseudo_thread_id`

System Variable	Name	`pseudo_thread_id`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	Type	integer

This variable is for internal server use.

- `query_alloc_block_size`

Command-Line Format	`--query-alloc-block-size=#`

System Variable	Name	query_alloc_block_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	8192
	Min Value	1024
	Max Value	4294967295
	Block Size	1024

The allocation size of memory blocks that are allocated for objects created during statement parsing and execution. If you have problems with memory fragmentation, it might help to increase this parameter.

- query_cache_limit

Command-Line Format	--query-cache-limit=#	
System Variable	Name	query_cache_limit
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	1048576
	Min Value	0
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	1048576
	Min Value	0
	Max Value	18446744073709551615

Do not cache results that are larger than this number of bytes. The default value is 1MB.

- query_cache_min_res_unit

Command-Line Format	--query-cache-min-res-unit=#	
System Variable	Name	query_cache_min_res_unit
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values (32-bit platforms)	**Type**	integer
	Default	`4096`
	Min Value	`512`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`4096`
	Min Value	`512`
	Max Value	`18446744073709551615`

The minimum size (in bytes) for blocks allocated by the query cache. The default value is 4096 (4KB). Tuning information for this variable is given in Section 8.10.3.3, “Query Cache Configuration”.

- `query_cache_size`

Command-Line Format	`--query-cache-size=#`	
System Variable	**Name**	`query_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms, 8.0.0)	**Type**	integer
	Default	`1048576`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (32-bit platforms, >= 8.0.1)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms, 8.0.0)	**Type**	integer
	Default	`1048576`
	Min Value	`0`
	Max Value	`18446744073709551615`
Permitted Values (64-bit platforms, >= 8.0.1)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709551615`

The amount of memory allocated for caching query results. By default, the query cache is disabled, achieved using defaults of `query_cache_size=0` and `query_cache_type=0`.

The permissible values are multiples of 1024; other values are rounded down to the nearest multiple. For nonzero values of `query_cache_size`, that many bytes of memory are allocated even if `query_cache_type=0`. See Section 8.10.3.3, “Query Cache Configuration”, for more information.

The query cache needs a minimum size of about 40KB to allocate its structures. (The exact size depends on system architecture.) If you set the value of `query_cache_size` too small, a warning will occur, as described in Section 8.10.3.3, “Query Cache Configuration”.

- `query_cache_type`

Command-Line Format	`--query-cache-type=#`	
System Variable	**Name**	`query_cache_type`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`0`
	Valid Values	`0`
		`1`
		`2`

Set the query cache type. Setting the `GLOBAL` value sets the type for all clients that connect thereafter. Individual clients can set the `SESSION` value to affect their own use of the query cache. Possible values are shown in the following table.

Option	**Description**
`0` or `OFF`	Do not cache results in or retrieve results from the query cache. Note that this does not deallocate the query cache buffer. To do that, you should set `query_cache_size` to 0.
`1` or `ON`	Cache all cacheable query results except for those that begin with `SELECT SQL_NO_CACHE`.
`2` or `DEMAND`	Cache results only for cacheable queries that begin with `SELECT SQL_CACHE`.

This variable defaults to `OFF`.

If the server is started with `query_cache_type` set to 0, it does not acquire the query cache mutex at all, which means that the query cache cannot be enabled at runtime and there is reduced overhead in query execution.

- `query_cache_wlock_invalidate`

Command-Line Format	`--query-cache-wlock-invalidate`	
System Variable	**Name**	`query_cache_wlock_invalidate`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean

	Default	FALSE

Normally, when one client acquires a WRITE lock on a MyISAM table, other clients are not blocked from issuing statements that read from the table if the query results are present in the query cache. Setting this variable to 1 causes acquisition of a WRITE lock for a table to invalidate any queries in the query cache that refer to the table. This forces other clients that attempt to access the table to wait while the lock is in effect.

- query_prealloc_size

Command-Line Format	--query-prealloc-size=#	
System Variable	**Name**	query_prealloc_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	8192
	Min Value	8192
	Max Value	4294967295
	Block Size	1024
Permitted Values (64-bit platforms)	**Type**	integer
	Default	8192
	Min Value	8192
	Max Value	18446744073709551615
	Block Size	1024

The size of the persistent buffer used for statement parsing and execution. This buffer is not freed between statements. If you are running complex queries, a larger query_prealloc_size value might be helpful in improving performance, because it can reduce the need for the server to perform memory allocation during query execution operations.

- rand_seed1

  The rand_seed1 and rand_seed2 variables exist as session variables only, and can be set but not read. The variables—but not their values—are shown in the output of SHOW VARIABLES.

  The purpose of these variables is to support replication of the RAND() function. For statements that invoke RAND(), the master passes two values to the slave, where they are used to seed the random number generator. The slave uses these values to set the session variables rand_seed1 and rand_seed2 so that RAND() on the slave generates the same value as on the master.

- rand_seed2

  See the description for rand_seed1.

- range_alloc_block_size

Command-Line Format	--range-alloc-block-size=#	
System Variable	Name	range_alloc_block_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	4096
	Min Value	4096
	Max Value	4294967295
	Block Size	1024
Permitted Values (64-bit platforms)	Type	integer
	Default	4096
	Min Value	4096
	Max Value	18446744073709547520
	Block Size	1024

The size of blocks that are allocated when doing range optimization.

- range_optimizer_max_mem_size

Command-Line Format	--range-optimizer-max-mem-size=N	
System Variable	Name	range_optimizer_max_mem_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	8388608
	Min Value	0
	Max Value	18446744073709551615

The limit on memory consumption for the range optimizer. A value of 0 means “no limit.” If an execution plan considered by the optimizer uses the range access method but the optimizer estimates that the amount of memory needed for this method would exceed the limit, it abandons the plan and considers other plans. For more information, see Limiting Memory Use for Range Optimization.

- rbr_exec_mode

System Variable	Name	rbr_exec_mode

	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`STRICT`
	Valid Values	`IDEMPOTENT`
		`STRICT`

This variable switches the server between `IDEMPOTENT` mode and `STRICT` mode. `IDEMPOTENT` mode causes suppression of duplicate-key and no-key-found errors. This mode is useful when replaying a row-based binary log on a server that causes conflicts with existing data. `mysqlbinlog` uses this mode when you set the `--idempotent` option by writing the following to the output:

```
SET SESSION RBR_EXEC_MODE=IDEMPOTENT;
```

- `read_buffer_size`

Command-Line Format	`--read-buffer-size=#`	
System Variable	**Name**	`read_buffer_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`131072`
	Min Value	`8200`
	Max Value	`2147479552`

Each thread that does a sequential scan for a `MyISAM` table allocates a buffer of this size (in bytes) for each table it scans. If you do many sequential scans, you might want to increase this value, which defaults to 131072. The value of this variable should be a multiple of 4KB. If it is set to a value that is not a multiple of 4KB, its value will be rounded down to the nearest multiple of 4KB.

This option is also used in the following context for all storage engines:

- For caching the indexes in a temporary file (not a temporary table), when sorting rows for `ORDER BY`.
- For bulk insert into partitions.
- For caching results of nested queries.

and in one other storage engine-specific way: to determine the memory block size for `MEMORY` tables.

The maximum permissible setting for `read_buffer_size` is 2GB.

For more information about memory use during different operations, see Section 8.12.3.1, “How MySQL Uses Memory”.

- `read_only`

Command-Line Format	--read-only	
System Variable	Name	read_only
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

When the read_only system variable is enabled, the server permits no client updates except from users who have the CONNECTION_ADMIN or SUPER privilege. This variable is disabled by default.

The server also supports a super_read_only system variable (disabled by default), which has these effects:

- If super_read_only is enabled, the server prohibits client updates, even from users who have the SUPER privilege.
- Setting super_read_only to ON implicitly forces read_only to ON.
- Setting read_only to OFF implicitly forces super_read_only to OFF.

Even with read_only enabled, the server permits these operations:

- Updates performed by slave threads, if the server is a replication slave. In replication setups, it can be useful to enable read_only on slave servers to ensure that slaves accept updates only from the master server and not from clients.
- Use of ANALYZE TABLE or OPTIMIZE TABLE statements. The purpose of read-only mode is to prevent changes to table structure or contents. Analysis and optimization do not qualify as such changes. This means, for example, that consistency checks on read-only replication slaves can be performed with mysqlcheck --all-databases --analyze.
- Operations on TEMPORARY tables.
- Inserts into the log tables (mysql.general_log and mysql.slow_log); see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”.
- Updates to Performance Schema tables, such as UPDATE or TRUNCATE TABLE operations.

Changes to read_only on a master server are not replicated to slave servers. The value can be set on a slave server independent of the setting on the master.

The following conditions apply to attempts to enable read_only (including implicit attempts resulting from enabling super_read_only):

- The attempt fails and an error occurs if you have any explicit locks (acquired with LOCK TABLES) or have a pending transaction.
- The attempt blocks while other clients hold explicit table locks or have pending transactions, until the locks are released and the transactions end. While the attempt to enable read_only is pending, requests by other clients for table locks or to begin transactions also block until read_only has been set.
- The attempt blocks if there are active transactions that hold metadata locks, until those transactions end.

  - `read_only` can be enabled while you hold a global read lock (acquired with `FLUSH TABLES WITH READ LOCK`) because that does not involve table locks.

- `read_rnd_buffer_size`

Command-Line Format	`--read-rnd-buffer-size=#`	
System Variable	Name	`read_rnd_buffer_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`262144`
	Min Value	`1`
	Max Value	`2147483647`

This variable is used for reads from `MyISAM` tables, and, for any storage engine, for Multi-Range Read optimization.

When reading rows from a `MyISAM` table in sorted order following a key-sorting operation, the rows are read through this buffer to avoid disk seeks. See Section 8.2.1.13, “ORDER BY Optimization”. Setting the variable to a large value can improve `ORDER BY` performance by a lot. However, this is a buffer allocated for each client, so you should not set the global variable to a large value. Instead, change the session variable only from within those clients that need to run large queries.

The maximum permissible setting for `read_rnd_buffer_size` is 2GB.

For more information about memory use during different operations, see Section 8.12.3.1, “How MySQL Uses Memory”. For information about Multi-Range Read optimization, see Section 8.2.1.10, “Multi-Range Read Optimization”.

- `relay_log_purge`

Command-Line Format	`--relay-log-purge`	
System Variable	Name	`relay_log_purge`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`TRUE`

Disables or enables automatic purging of relay log files as soon as they are not needed any more. The default value is 1 (`ON`).

- `relay_log_space_limit`

Command-Line Format	`--relay-log-space-limit=#`	
System Variable	Name	`relay_log_space_limit`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709551615`

The maximum amount of space to use for all relay logs.

- `report_host`

Command-Line Format	`--report-host=host_name`	
System Variable	**Name**	`report_host`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string

The value of the `--report-host` option.

- `report_password`

Command-Line Format	`--report-password=name`	
System Variable	**Name**	`report_password`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string

The value of the `--report-password` option. Not the same as the password used for the MySQL replication user account.

- `report_port`

Command-Line Format	`--report-port=#`	
System Variable	**Name**	`report_port`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer

	Default	`[slave_port]`
	Min Value	`0`
	Max Value	`65535`

The value of the `--report-port` option.

- `report_user`

Command-Line Format	`--report-user=name`	
System Variable	**Name**	`report_user`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string

The value of the `--report-user` option. Not the same as the name for the MySQL replication user account.

- `require_secure_transport`

Command-Line Format	`--require-secure-transport[={OFF|ON}]`	
System Variable	**Name**	`require_secure_transport`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Whether client connections to the server are required to use some form of secure transport. When this variable is enabled, the server permits only TCP/IP connections that use SSL, or connections that use a socket file (on Unix) or shared memory (on Windows). The server rejects nonsecure connection attempts, which fail with an `ER_SECURE_TRANSPORT_REQUIRED` error.

This capability supplements per-account SSL requirements, which take precedence. For example, if an account is defined with `REQUIRE SSL`, enabling `require_secure_transport` does not make it possible to use the account to connect using a Unix socket file.

It is possible for a server to have no secure transports available. For example, a server on Windows supports no secure transports if started without specifying any SSL certificate or key files and with the `shared_memory` system variable disabled. Under these conditions, attempts to enable `require_secure_transport` at startup cause the server to write a message to the error log and exit. Attempts to enable the variable at runtime fail with an `ER_NO_SECURE_TRANSPORTS_CONFIGURED` error.

- `rpl_semi_sync_master_enabled`

System Variable	**Name**	`rpl_semi_sync_master_enabled`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Controls whether semisynchronous replication is enabled on the master. To enable or disable the plugin, set this variable to `ON` or `OFF` (or 1 or 0), respectively. The default is `OFF`.

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_master_timeout`

System Variable	**Name**	`rpl_semi_sync_master_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`10000`

A value in milliseconds that controls how long the master waits on a commit for acknowledgment from a slave before timing out and reverting to asynchronous replication. The default value is 10000 (10 seconds).

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_master_trace_level`

System Variable	**Name**	`rpl_semi_sync_master_trace_level`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`32`

The semisynchronous replication debug trace level on the master. Four levels are defined:

- 1 = general level (for example, time function failures)
- 16 = detail level (more verbose information)
- 32 = net wait level (more information about network waits)
- 64 = function level (information about function entry and exit)

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_master_wait_for_slave_count`

System Variable	**Name**	`rpl_semi_sync_master_wait_for_slave_count`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	integer
	Default	`1`
	Min Value	`1`
	Max Value	`65535`

The number of slave acknowledgments the master must receive per transaction before proceeding. By default `rpl_semi_sync_master_wait_for_slave_count` is `1`, meaning that semisynchronous replication proceeds after receiving a single slave acknowledgment. Performance is best for small values of this variable.

For example, if `rpl_semi_sync_master_wait_for_slave_count` is `2`, then 2 slaves must acknowledge receipt of the transaction before the timeout period configured by `rpl_semi_sync_master_timeout` for semisynchronous replication to proceed. If less slaves acknowledge receipt of the transaction during the timeout period, the master reverts to normal replication.

Note

This behavior also depends on `rpl_semi_sync_master_wait_no_slave`

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_master_wait_no_slave`

System Variable	**Name**	`rpl_semi_sync_master_wait_no_slave`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Controls whether the master waits for the timeout period configured by `rpl_semi_sync_master_timeout` to expire, even if the slave count drops to less than the number of slaves configured by `rpl_semi_sync_master_wait_for_slave_count` during the timeout period.

When the value of `rpl_semi_sync_master_wait_no_slave` is `ON` (the default), it is permissible for the slave count to drop to less than `rpl_semi_sync_master_wait_for_slave_count` during the timeout period. As long as enough slaves acknowledge the transaction before the timeout period expires, semisynchronous replication continues.

When the value of `rpl_semi_sync_master_wait_no_slave` is `OFF`, if the slave count drops to less than the number configured in `rpl_semi_sync_master_wait_for_slave_count` at any time during the timeout period configured by `rpl_semi_sync_master_timeout`, the master reverts to normal replication.

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_master_wait_point`

System Variable	**Name**	`rpl_semi_sync_master_wait_point`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`AFTER_SYNC`
	Valid Values	`AFTER_SYNC`
		`AFTER_COMMIT`

This variable controls the point at which a semisynchronous replication master waits for slave acknowledgment of transaction receipt before returning a status to the client that committed the transaction. These values are permitted:

- `AFTER_SYNC` (the default): The master writes each transaction to its binary log and the slave, and syncs the binary log to disk. The master waits for slave acknowledgment of transaction receipt after the sync. Upon receiving acknowledgment, the master commits the transaction to the storage engine and returns a result to the client, which then can proceed.

- `AFTER_COMMIT`: The master writes each transaction to its binary log and the slave, syncs the binary log, and commits the transaction to the storage engine. The master waits for slave acknowledgment of transaction receipt after the commit. Upon receiving acknowledgment, the master returns a result to the client, which then can proceed.

The replication characteristics of these settings differ as follows:

- With `AFTER_SYNC`, all clients see the committed transaction at the same time: After it has been acknowledged by the slave and committed to the storage engine on the master. Thus, all clients see the same data on the master.

  In the event of master failure, all transactions committed on the master have been replicated to the slave (saved to its relay log). A crash of the master and failover to the slave is lossless because the slave is up to date.

- With `AFTER_COMMIT`, the client issuing the transaction gets a return status only after the server commits to the storage engine and receives slave acknowledgment. After the commit and before slave acknowledgment, other clients can see the committed transaction before the committing client.

  If something goes wrong such that the slave does not process the transaction, then in the event of a master crash and failover to the slave, it is possible that such clients will see a loss of data relative to what they saw on the master.

This variable is available only if the master-side semisynchronous replication plugin is installed.

With the addition of `rpl_semi_sync_master_wait_point` in MySQL 5.7, a version compatibility constraint was created because it increments the semisynchronous interface version: Servers for MySQL 5.7 and higher do not work with semisynchronous replication plugins from older versions, nor do servers from older versions work with semisynchronous replication plugins for MySQL 5.7 and higher.

- `rpl_semi_sync_slave_enabled`

System Variable	**Name**	`rpl_semi_sync_slave_enabled`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Controls whether semisynchronous replication is enabled on the slave. To enable or disable the plugin, set this variable to `ON` or `OFF` (or 1 or 0), respectively. The default is `OFF`.

This variable is available only if the slave-side semisynchronous replication plugin is installed.

- `rpl_semi_sync_slave_trace_level`

System Variable	**Name**	`rpl_semi_sync_slave_trace_level`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`32`

The semisynchronous replication debug trace level on the slave. See `rpl_semi_sync_master_trace_level` for the permissible values.

This variable is available only if the slave-side semisynchronous replication plugin is installed.

- `schema_definition_cache`

Command-Line Format	`--schema-definition-cache=N`	
System Variable	**Name**	`schema_definition_cache`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`256`
	Min Value	`256`
	Max Value	`524288`

Defines a limit for the number of schema definition objects, both used and unused, that can be kept in the dictionary object cache.

Unused schema definition objects are only kept in the dictionary object cache when the number in use is less than the capacity defined by `schema_definition_cache`.

A setting of `0` means that schema definition objects are only kept in the dictionary object cache while they are in use.

For more information, see Section 14.4, “Dictionary Object Cache”.

- `secure_auth`

Command-Line Format	--secure-auth	
System Variable	Name	secure_auth
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	ON
	Valid Values	ON

If this variable is enabled, the server blocks connections by clients that attempt to use accounts that have passwords stored in the old (pre-4.1) format.

Enable this variable to prevent all use of passwords employing the old format (and hence insecure communication over the network).

This variable is deprecated and will be removed in a future MySQL release. It is always enabled and attempting to disable it produces an error.

Server startup fails with an error if this variable is enabled and the privilege tables are in pre-4.1 format.

Note

Passwords that use the pre-4.1 hashing method are less secure than passwords that use the native password hashing method and should be avoided. Pre-4.1 passwords are deprecated and support for them has been removed.

- secure_file_priv

Command-Line Format	--secure-file-priv=dir_name	
System Variable	Name	secure_file_priv
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	platform specific
	Valid Values	empty
		dirname
		NULL

This variable is used to limit the effect of data import and export operations, such as those performed by the LOAD DATA and SELECT ... INTO OUTFILE statements and the LOAD_FILE() function. These operations are permitted only to users who have the FILE privilege.

secure_file_priv may be set as follows:

- If empty, the variable has no effect. This is not a secure setting.
- If set to the name of a directory, the server limits import and export operations to work only with files in that directory. The directory must exist; the server will not create it.

- If set to `NULL`, the server disables import and export operations.

The default value is platform specific and depends on the value of the `INSTALL_LAYOUT` `CMake` option, as shown in the following table. To specify the default `secure_file_priv` value explicitly if you are building from source, use the `INSTALL_SECURE_FILE_PRIVDIR` `CMake` option.

`INSTALL_LAYOUT` Value	**Default `secure_file_priv` Value**
`STANDALONE`, `WIN`	empty
`DEB`, `RPM`, `SLES`, `SVR4`	`/var/lib/mysql-files`
Otherwise	`mysql-files` under the `CMAKE_INSTALL_PREFIX` value

The server checks the value of `secure_file_priv` at startup and writes a warning to the error log if the value is insecure. A non-`NULL` value is considered insecure if it is empty, or the value is the data directory or a subdirectory of it, or a directory that is accessible by all users. If `secure_file_priv` is set to a nonexistent path, the server writes an error message to the error log and exits.

- `server_id`

Command-Line Format	`--server-id=#`	
System Variable	**Name**	`server_id`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

The server ID, used in replication to give each master and slave a unique identity. This variable is set by the `--server-id` option. For each server participating in replication, you should pick a positive integer in the range from 1 to 2^{32} − 1 to act as that server's ID.

- `session_track_gtids`

Command-Line Format	`--session-track-gtids=[value]`	
System Variable	**Name**	`session_track_gtids`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`OFF`
	Valid Values	`OFF`
		`OWN_GTID`
		`ALL_GTIDS`

Controls a tracker for capturing GTIDs and returning them in the OK packet. Depending on the value of this option, at the end of executing a transaction, the GTIDs specified are captured by the tracker and appended to the OK packet. The possible sets of GTIDs to track are:

- `OFF` means that no GTIDs are included in the OK packet. This is the same behavior as versions of MySQL prior to 5.7.6.
- `OWN_GTID` configures the tracker to collect GTIDs generated by successfully committed read/write transactions.
- `ALL_GTIDS` configures the tracker to collect all of the GTIDs in `gtid_executed` at the time the current transaction commits, regardless of whether the transaction is read/write or read-only.

For information about obtaining session state-change information within client programs, see Section 27.7.7.67, “mysql_session_track_get_first()”.

- `session_track_schema`

Command-Line Format	`--session-track-schema=#`	
System Variable	**Name**	`session_track_schema`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Controls whether the server tracks changes to the default schema (database) name within the current session and makes this information available to the client when changes occur.

If notification is enabled, any setting of the default schema is reported, even if the new schema name is the same as the old.

For information about obtaining session state-change information within client programs, see Section 27.7.7.67, “mysql_session_track_get_first()”.

- `session_track_state_change`

Command-Line Format	`--session-track-state-change=#`	
System Variable	**Name**	`session_track_state_change`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Controls whether the server tracks changes to the state of the current session and notifies the client when state changes occur. Session state consists of these values:

- The default schema (database)
- Session-specific values for system variables
- User-defined variables

- Temporary tables
- Prepared statements

If the session-state tracker is enabled, any assignments to session state values are reported, even if the new values are the same as the old.

The `session_track_state_change` variable controls only notification of when changes occur, not what the changes are. To receive notification for changes to the default schema name and session system variable values, use the `session_track_schema` and `session_track_system_variables` system variables.

For information about obtaining session state-change information within client programs, see Section 27.7.7.67, “mysql_session_track_get_first()”.

- `session_track_system_variables`

Command-Line Format	`--session-track-system-variables=#`	
System Variable	**Name**	`session_track_system_variables`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`time_zone, autocommit, character_set_client, character_set_results, character_set_connection`

Controls whether the server tracks changes to the session system variables and makes this information available to the client when changes occur. The variable value is a comma-separated list of variables for which to track changes. By default, notification is enabled for `time_zone`, `autocommit`, `character_set_client`, `character_set_results`, and `character_set_connection`. (The latter three variables are those affected by `SET NAMES`.)

The special value `*` causes the server to track changes to all session variables. If given, this value must be specified by itself without specific system variable names.

Notification occurs for all assignments to tracked session system variables, even if the new values are the same as the old.

For information about obtaining session state-change information within client programs, see Section 27.7.7.67, “mysql_session_track_get_first()”.

- `sha256_password_auto_generate_rsa_keys`

Command-Line Format	`--sha256-password-auto-generate-rsa-keys[={OFF|ON}]`	
System Variable	**Name**	`sha256_password_auto_generate_rsa_keys`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`ON`

This variable is available if the server was compiled using OpenSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). It controls whether the server autogenerates RSA private/public key-pair files in the data directory, if they do not already exist.

At startup, the server automatically generates RSA private/public key-pair files in the data directory if the `sha256_password_auto_generate_rsa_keys` system variable is enabled, no RSA options are specified, and the RSA files are missing from the data directory. These files enable secure password exchange using RSA over unencrypted connections for accounts authenticated by the `sha256_password` plugin; see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

For more information about RSA file autogeneration, including file names and characteristics, see Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”

The `auto_generate_certs` system variable is related but controls autogeneration of SSL certificate and key files needed for secure connections using SSL.

- `sha256_password_private_key_path`

System Variable	Name	sha256_password_private_key_path
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name
	Default	private_key.pem

This variable is available if MySQL was compiled using OpenSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). Its value is the path name of the RSA private key file for the `sha256_password` authentication plugin. If the file is named as a relative path, it is interpreted relative to the server data directory. The file must be in PEM format. Because this file stores a private key, its access mode should be restricted so that only the MySQL server can read it.

For information about `sha256_password`, including instructions for creating the RSA key files, see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

- `sha256_password_proxy_users`

Command-Line Format	--sha256-password-proxy-users=[={OFF|ON}]	
System Variable	Name	sha256_password_proxy_users
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

This variable controls whether the `sha256_password` built-in authentication plugin supports proxy users. It has no effect unless the `check_proxy_users` system variable is enabled. For information about user proxying, see Section 6.3.11, “Proxy Users”.

- `sha256_password_public_key_path`

System Variable	Name	sha256_password_public_key_path

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name
	Default	`public_key.pem`

This variable is available if MySQL was compiled using OpenSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). Its value is the path name of the RSA public key file for the `sha256_password` authentication plugin. If the file is named as a relative path, it is interpreted relative to the server data directory. The file must be in PEM format. Because this file stores a public key, copies can be freely distributed to client users. (Clients that explicitly specify a public key when connecting to the server using RSA password encryption must use the same public key as that used by the server.)

For information about `sha256_password`, including instructions for creating the RSA key files and how clients specify the RSA public key, see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

- `shared_memory`

Command-Line Format	`--shared-memory[={0,1}]`	
System Variable	**Name**	`shared_memory`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Windows	
Permitted Values	**Type**	boolean
	Default	`FALSE`

(Windows only.) Whether the server permits shared-memory connections.

- `shared_memory_base_name`

Command-Line Format	`--shared-memory-base-name=name`	
System Variable	**Name**	`shared_memory_base_name`
	Variable Scope	Global
	Dynamic Variable	No
Platform Specific	Windows	
Permitted Values	**Type**	string
	Default	`MYSQL`

(Windows only.) The name of shared memory to use for shared-memory connections. This is useful when running multiple MySQL instances on a single physical machine. The default name is `MYSQL`. The name is case sensitive.

- `show_compatibility_56`

Deprecated	5.7.6
Removed	8.0.1
Command-Line Format	`--show-compatibility-56[={OFF|ON}]`

System Variable	Name	show_compatibility_56
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

This variable was used in the transition period during which system and status variable information in INFORMATION_SCHEMA tables was moved to Performance Schema tables. That transition period ended in MySQL 8.0.1, at which time this variable was removed. For advice on migrating away from the INFORMATION_SCHEMA tables to the Performance Schema tables, see Migrating to Performance Schema System and Status Variable Tables.

- show_old_temporals

Deprecated	5.7.6	
Command-Line Format	--show-old-temporals={OFF|ON}	
System Variable	Name	show_old_temporals
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Whether SHOW CREATE TABLE output includes comments to flag temporal columns found to be in pre-5.6.4 format (TIME, DATETIME, and TIMESTAMP columns without support for fractional seconds precision). This variable is disabled by default. If enabled, SHOW CREATE TABLE output looks like this:

```
CREATE TABLE `mytbl` (
  `ts` timestamp /* 5.5 binary format */ NOT NULL DEFAULT CURRENT_TIMESTAMP,
  `dt` datetime /* 5.5 binary format */ DEFAULT NULL,
  `t` time /* 5.5 binary format */ DEFAULT NULL
) DEFAULT CHARSET=utf8mb4
```

Output for the COLUMN_TYPE column of the INFORMATION_SCHEMA.COLUMNS table is affected similarly.

This variable is deprecated and will be removed in a future MySQL release.

- skip_external_locking

Command-Line Format	--skip-external-locking	
System Variable	Name	skip_external_locking
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	ON

This is `OFF` if `mysqld` uses external locking (system locking), `ON` if external locking is disabled. This affects only `MyISAM` table access.

This variable is set by the `--external-locking` or `--skip-external-locking` option. External locking is disabled by default.

External locking affects only `MyISAM` table access. For more information, including conditions under which it can and cannot be used, see Section 8.11.5, “External Locking”.

- `skip_name_resolve`

Command-Line Format	`--skip-name-resolve`	
System Variable	Name	`skip_name_resolve`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`OFF`

This variable is set from the value of the `--skip-name-resolve` option. If it is `OFF`, `mysqld` resolves host names when checking client connections. If it is `ON`, `mysqld` uses only IP numbers; in this case, all `Host` column values in the grant tables must be IP addresses or `localhost`. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.

- `skip_networking`

Command-Line Format	`--skip-networking`	
System Variable	Name	`skip_networking`
	Variable Scope	Global
	Dynamic Variable	No

This is `ON` if the server permits only local (non-TCP/IP) connections. On Unix, local connections use a Unix socket file. On Windows, local connections use a named pipe or shared memory. This variable can be set to `ON` with the `--skip-networking` option.

- `skip_show_database`

Command-Line Format	`--skip-show-database`	
System Variable	Name	`skip_show_database`
	Variable Scope	Global
	Dynamic Variable	No

This prevents people from using the `SHOW DATABASES` statement if they do not have the `SHOW DATABASES` privilege. This can improve security if you have concerns about users being able to see databases belonging to other users. Its effect depends on the `SHOW DATABASES` privilege: If the variable value is `ON`, the `SHOW DATABASES` statement is permitted only to users who have the `SHOW DATABASES` privilege, and the statement displays all database names. If the value is `OFF`, `SHOW DATABASES` is permitted to all users, but displays the names of only those databases for which the

user has the `SHOW DATABASES` or other privilege. (Note that *any* global privilege is considered a privilege for the database.)

- `slow_launch_time`

Command-Line Format	`--slow-launch-time=#`	
System Variable	Name	`slow_launch_time`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`2`

If creating a thread takes longer than this many seconds, the server increments the `Slow_launch_threads` status variable.

- `slow_query_log`

Command-Line Format	`--slow-query-log`	
System Variable	Name	`slow_query_log`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Whether the slow query log is enabled. The value can be 0 (or `OFF`) to disable the log or 1 (or `ON`) to enable the log. The default value depends on whether the `--slow_query_log` option is given. The destination for log output is controlled by the `log_output` system variable; if that value is `NONE`, no log entries are written even if the log is enabled.

"Slow" is determined by the value of the `long_query_time` variable. See Section 5.4.5, "The Slow Query Log".

- `slow_query_log_file`

Command-Line Format	`--slow-query-log-file=file_name`	
System Variable	Name	`slow_query_log_file`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	file name
	Default	`host_name-slow.log`

The name of the slow query log file. The default value is *`host_name`*`-slow.log`, but the initial value can be changed with the `--slow_query_log_file` option.

- `socket`

Command-Line Format	`--socket={file_name|pipe_name}`

System Variable	Name	socket
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	/tmp/mysql.sock

On Unix platforms, this variable is the name of the socket file that is used for local client connections. The default is /tmp/mysql.sock. (For some distribution formats, the directory might be different, such as /var/lib/mysql for RPMs.)

On Windows, this variable is the name of the named pipe that is used for local client connections. The default value is MySQL (not case sensitive).

- sort_buffer_size

Command-Line Format	--sort-buffer-size=#	
System Variable	Name	sort_buffer_size
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (Windows)	Type	integer
	Default	262144
	Min Value	32768
	Max Value	4294967295
Permitted Values (Other, 32-bit platforms)	Type	integer
	Default	262144
	Min Value	32768
	Max Value	4294967295
Permitted Values (Other, 64-bit platforms)	Type	integer
	Default	262144
	Min Value	32768
	Max Value	18446744073709551615

Each session that must perform a sort allocates a buffer of this size. sort_buffer_size is not specific to any storage engine and applies in a general manner for optimization. At minimum the sort_buffer_size value must be large enough to accommodate fifteen tuples in the sort buffer. Also, increasing the value of max_sort_length may require increasing the value of sort_buffer_size. For more information, see Section 8.2.1.13, “ORDER BY Optimization”

If you see many Sort_merge_passes per second in SHOW GLOBAL STATUS output, you can consider increasing the sort_buffer_size value to speed up ORDER BY or GROUP BY operations that cannot be improved with query optimization or improved indexing.

The optimizer tries to work out how much space is needed but can allocate more, up to the limit. Setting it larger than required globally will slow down most queries that sort. It is best to increase it as a session setting, and only for the sessions that need a larger size. On Linux, there are thresholds of 256KB and 2MB where larger values may significantly slow down memory allocation, so you should consider staying below one of those values. Experiment to find the best value for your workload. See Section B.5.3.5, “Where MySQL Stores Temporary Files”.

The maximum permissible setting for `sort_buffer_size` is 4GB−1. Larger values are permitted for 64-bit platforms (except 64-bit Windows, for which large values are truncated to 4GB−1 with a warning).

- `sql_auto_is_null`

System Variable	Name	`sql_auto_is_null`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`0`

If this variable is set to 1, then after a statement that successfully inserts an automatically generated `AUTO_INCREMENT` value, you can find that value by issuing a statement of the following form:

```
SELECT * FROM tbl_name WHERE auto_col IS NULL
```

If the statement returns a row, the value returned is the same as if you invoked the `LAST_INSERT_ID()` function. For details, including the return value after a multiple-row insert, see Section 12.14, “Information Functions”. If no `AUTO_INCREMENT` value was successfully inserted, the `SELECT` statement returns no row.

The behavior of retrieving an `AUTO_INCREMENT` value by using an `IS NULL` comparison is used by some ODBC programs, such as Access. See Obtaining Auto-Increment Values. This behavior can be disabled by setting `sql_auto_is_null` to 0.

The default value of `sql_auto_is_null` is 0.

- `sql_big_selects`

System Variable	Name	`sql_big_selects`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`1`

If set to 0, MySQL aborts `SELECT` statements that are likely to take a very long time to execute (that is, statements for which the optimizer estimates that the number of examined rows exceeds the value of `max_join_size`). This is useful when an inadvisable `WHERE` statement has been issued. The default value for a new connection is 1, which permits all `SELECT` statements.

If you set the `max_join_size` system variable to a value other than `DEFAULT`, `sql_big_selects` is set to 0.

- `sql_buffer_result`

System Variable	Name	sql_buffer_result
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	0

If set to 1, `sql_buffer_result` forces results from `SELECT` statements to be put into temporary tables. This helps MySQL free the table locks early and can be beneficial in cases where it takes a long time to send results to the client. The default value is 0.

- `sql_log_bin`

System Variable	Name	sql_log_bin
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean

This variable controls whether logging to the binary log is done. The default value is 1 (do logging). To change logging for the current session, change the session value of this variable. The session user must have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege to set this variable.

Setting this variable to 0 prevents GTIDs from being assigned to transactions in the binary log. If you are using GTIDs for replication, this means that, even when binary logging is later enabled once again, the GTIDs written into the log from this point do not account for any transactions that occurred in the meantime—in effect, those transactions are lost.

In MySQL 8.0, it is not possible to set `@@session.sql_log_bin` within a transaction or subquery. (Bug #53437)

- `sql_log_off`

System Variable	Name	sql_log_off
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	0

This variable controls whether logging to the general query log is done. The default value is 0 (do logging). To change logging for the current session, change the session value of this variable. The session user must have the `SUPER` privilege to set this option. The default value is 0.

- `sql_mode`

Command-Line Format	--sql-mode=name	
System Variable	Name	sql_mode

	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	set
	Default	`ONLY_FULL_GROUP_BY STRICT_TRANS_TABLES NO_ZERO_IN_DATE NO_ZERO_DATE ERROR_FOR_DIVISION_BY_ZERO NO_AUTO_CREATE_USER NO_ENGINE_SUBSTITUTION`
	Valid Values	`ALLOW_INVALID_DATES`
		`ANSI_QUOTES`
		`ERROR_FOR_DIVISION_BY_ZERO`
		`HIGH_NOT_PRECEDENCE`
		`IGNORE_SPACE`
		`NO_AUTO_CREATE_USER`
		`NO_AUTO_VALUE_ON_ZERO`
		`NO_BACKSLASH_ESCAPES`
		`NO_DIR_IN_CREATE`
		`NO_ENGINE_SUBSTITUTION`
		`NO_FIELD_OPTIONS`
		`NO_KEY_OPTIONS`
		`NO_TABLE_OPTIONS`
		`NO_UNSIGNED_SUBTRACTION`
		`NO_ZERO_DATE`
		`NO_ZERO_IN_DATE`
		`ONLY_FULL_GROUP_BY`
		`PAD_CHAR_TO_FULL_LENGTH`
		`PIPES_AS_CONCAT`
		`REAL_AS_FLOAT`
		`STRICT_ALL_TABLES`
		`STRICT_TRANS_TABLES`
Permitted Values (>= 8.0.1)	**Type**	set
	Default	`ONLY_FULL_GROUP_BY STRICT_TRANS_TABLES NO_ZERO_IN_DATE NO_ZERO_DATE ERROR_FOR_DIVISION_BY_ZERO NO_AUTO_CREATE_USER NO_ENGINE_SUBSTITUTION`
	Valid Values	`ALLOW_INVALID_DATES`
		`ANSI_QUOTES`
		`ERROR_FOR_DIVISION_BY_ZERO`
		`HIGH_NOT_PRECEDENCE`
		`IGNORE_SPACE`
		`NO_AUTO_CREATE_USER`
		`NO_AUTO_VALUE_ON_ZERO`
		`NO_BACKSLASH_ESCAPES`

		NO_DIR_IN_CREATE
		NO_ENGINE_SUBSTITUTION
		NO_FIELD_OPTIONS
		NO_KEY_OPTIONS
		NO_TABLE_OPTIONS
		NO_UNSIGNED_SUBTRACTION
		NO_ZERO_DATE
		NO_ZERO_IN_DATE
		ONLY_FULL_GROUP_BY
		PAD_CHAR_TO_FULL_LENGTH
		PIPES_AS_CONCAT
		REAL_AS_FLOAT
		STRICT_ALL_TABLES
		STRICT_TRANS_TABLES
		TIME_TRUNCATE_FRACTIONAL

The current server SQL mode, which can be set dynamically. For details, see Section 5.1.8, “Server SQL Modes”.

Note

MySQL installation programs may configure the SQL mode during the installation process.

If the SQL mode differs from the default or from what you expect, check for a setting in an option file that the server reads at startup.

- `sql_notes`

  If set to 1 (the default), warnings of `Note` level increment `warning_count` and the server records them. If set to 0, `Note` warnings do not increment `warning_count` and the server does not record them. `mysqldump` includes output to set this variable to 0 so that reloading the dump file does not produce warnings for events that do not affect the integrity of the reload operation.

- `sql_quote_show_create`

  If set to 1 (the default), the server quotes identifiers for `SHOW CREATE TABLE` and `SHOW CREATE DATABASE` statements. If set to 0, quoting is disabled. This option is enabled by default so that replication works for identifiers that require quoting. See Section 13.7.5.10, “SHOW CREATE TABLE Syntax”, and Section 13.7.5.6, “SHOW CREATE DATABASE Syntax”.

- `sql_safe_updates`

  If set to 1, MySQL aborts `UPDATE` or `DELETE` statements that do not use a key in the `WHERE` clause or a `LIMIT` clause. (Specifically, `UPDATE` statements must have a `WHERE` clause that uses a key or a `LIMIT` clause, or both. `DELETE` statements must have both.) This makes it possible to catch `UPDATE` or `DELETE` statements where keys are not used properly and that would probably change or delete a large number of rows. The default value is 0.

- `sql_select_limit`

System Variable	**Name**	`sql_select_limit`
	Variable Scope	Global, Session

	Dynamic Variable	Yes
Permitted Values	**Type**	integer

The maximum number of rows to return from `SELECT` statements. The default value for a new connection is the maximum number of rows that the server permits per table. Typical default values are $(2^{32})-1$ or $(2^{64})-1$. If you have changed the limit, the default value can be restored by assigning a value of `DEFAULT`.

If a `SELECT` has a `LIMIT` clause, the `LIMIT` takes precedence over the value of `sql_select_limit`.

- `sql_warnings`

  This variable controls whether single-row `INSERT` statements produce an information string if warnings occur. The default is 0. Set the value to 1 to produce an information string.

- `ssl_ca`

Command-Line Format	`--ssl-ca=file_name`	
System Variable	**Name**	`ssl_ca`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The path to a file with a list of trusted SSL CAs.

- `ssl_capath`

Command-Line Format	`--ssl-capath=dir_name`	
System Variable	**Name**	`ssl_capath`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The path to a directory that contains trusted SSL CA certificates in PEM format.

- `ssl_cert`

Command-Line Format	`--ssl-cert=file_name`	
System Variable	**Name**	`ssl_cert`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The name of the SSL certificate file to use for establishing a secure connection.

- `ssl_cipher`

Command-Line Format	--ssl-cipher=name	
System Variable	Name	ssl_cipher
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

A list of permissible ciphers to use for SSL encryption.

- ssl_crl

Command-Line Format	--ssl-crl=file_name	
System Variable	Name	ssl_crl
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The path to a file containing certificate revocation lists in PEM format. Revocation lists work for MySQL distributions compiled using OpenSSL (but not yaSSL). See Section 6.4.1, “OpenSSL Versus yaSSL”.

- ssl_crlpath

Command-Line Format	--ssl-crlpath=dir_name	
System Variable	Name	ssl_crlpath
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The path to a directory that contains files containing certificate revocation lists in PEM format. Revocation lists work for MySQL distributions compiled using OpenSSL (but not yaSSL). See Section 6.4.1, “OpenSSL Versus yaSSL”.

- ssl_key

Command-Line Format	--ssl-key=file_name	
System Variable	Name	ssl_key
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The name of the SSL key file to use for establishing a secure connection.

Command-Line Format	--stored-program-cache=#	
System Variable	Name	stored_program_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	256
	Min Value	16
	Max Value	524288

Sets a soft upper limit for the number of cached stored routines per connection. The value of this variable is specified in terms of the number of stored routines held in each of the two caches maintained by the MySQL Server for, respectively, stored procedures and stored functions.

Whenever a stored routine is executed this cache size is checked before the first or top-level statement in the routine is parsed; if the number of routines of the same type (stored procedures or stored functions according to which is being executed) exceeds the limit specified by this variable, the corresponding cache is flushed and memory previously allocated for cached objects is freed. This allows the cache to be flushed safely, even when there are dependencies between stored routines.

The stored procedure and stored function caches exists in parallel with the stored program definition cache partition of the dictionary object cache. The stored procedure and stored function caches are per connection, while the stored program definition cache is shared. The existence of objects in the stored procedure and stored function caches have no dependence on the existence of objects in the stored program definition cache, and vice versa. For more information, see Section 14.4, “Dictionary Object Cache”.

- stored_program_definition_cache

Command-Line Format	--stored-program-definition-cache=N	
System Variable	Name	stored_program_definition_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	256
	Min Value	256
	Max Value	524288

Defines a limit for the number of stored program definition objects, both used and unused, that can be kept in the dictionary object cache.

Unused stored program definition objects are only kept in the dictionary object cache when the number in use is less than the capacity defined by stored_program_definition_cache.

A setting of 0 means that stored program definition objects are only kept in the dictionary object cache while they are in use.

The stored program definition cache partition exists in parallel with the stored procedure and stored function caches that are configured using the `stored_program_cache` option.

The `stored_program_cache` option sets a soft upper limit for the number of cached stored procedures or functions per connection, and the limit is checked each time a connection executes a stored procedure or function. The stored program definition cache partition, on the other hand, is a shared cache that stores stored program definition objects for other purposes. The existence of objects in the stored program definition cache partition has no dependence on the existence of objects in the stored procedure cache or stored function cache, and vice versa.

For related information, see Section 14.4, “Dictionary Object Cache”.

- `super_read_only`

Command-Line Format	`--super-read-only[={OFF|ON}]`	
System Variable	**Name**	`super_read_only`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

If the `read_only` system variable is enabled, the server permits client updates only from users who have the `SUPER` privilege. If the `super_read_only` system variable is also enabled, the server prohibits client updates even from users who have `SUPER`. See the description of the `read_only` system variable for a description of read-only mode and information about how `read_only` and `super_read_only` interact.

Changes to `super_read_only` on a master server are not replicated to slave servers. The value can be set on a slave server independent of the setting on the master.

- `system_time_zone`

System Variable	**Name**	`system_time_zone`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string

The server system time zone. When the server begins executing, it inherits a time zone setting from the machine defaults, possibly modified by the environment of the account used for running the server or the startup script. The value is used to set `system_time_zone`. Typically the time zone is specified by the `TZ` environment variable. It also can be specified using the `--timezone` option of the `mysqld_safe` script.

The `system_time_zone` variable differs from `time_zone`. Although they might have the same value, the latter variable is used to initialize the time zone for each client that connects. See Section 10.6, “MySQL Server Time Zone Support”.

- `tablespace_definition_cache`

Command-Line Format	--tablespace-definition-cache=N	
System Variable	Name	tablespace_definition_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	256
	Min Value	256
	Max Value	524288

Defines a limit for the number of tablespace definition objects, both used and unused, that can be kept in the dictionary object cache.

Unused tablespace definition objects are only kept in the dictionary object cache when the number in use is less than the capacity defined by tablespace_definition_cache.

A setting of 0 means that tablespace definition objects are only kept in the dictionary object cache while they are in use.

For more information, see Section 14.4, “Dictionary Object Cache”.

- table_definition_cache

System Variable	Name	table_definition_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	-1 (autosized)
	Min Value	400
	Max Value	524288

The number of table definitions that can be stored in the definition cache. If you use a large number of tables, you can create a large table definition cache to speed up opening of tables. The table definition cache takes less space and does not use file descriptors, unlike the normal table cache. The minimum value is 400. The default value is based on the following formula, capped to a limit of 2000:

```
400 + (table_open_cache / 2)
```

For InnoDB, table_definition_cache acts as a soft limit for the number of open table instances in the InnoDB data dictionary cache. If the number of open table instances exceeds the table_definition_cache setting, the LRU mechanism begins to mark table instances for eviction and eventually removes them from the data dictionary cache. The limit helps address situations in which significant amounts of memory would be used to cache rarely used table instances until the next server restart. The number of table instances with cached metadata could

be higher than the limit defined by `table_definition_cache`, because `InnoDB` system table instances and parent and child table instances with foreign key relationships are not placed on the LRU list and are not subject to eviction from memory.

Additionally, `table_definition_cache` defines a soft limit for the number of `InnoDB` file-per-table tablespaces that can be open at one time, which is also controlled by `innodb_open_files`. If both `table_definition_cache` and `innodb_open_files` are set, the highest setting is used. If neither variable is set, `table_definition_cache`, which has a higher default value, is used. If the number of open tablespace file handles exceeds the limit defined by `table_definition_cache` or `innodb_open_files`, the LRU mechanism searches the tablespace file LRU list for files that are fully flushed and are not currently being extended. This process is performed each time a new tablespace is opened. If there are no “inactive” tablespaces, no tablespace files are closed.

The table definition cache exists in parallel with the table definition cache partition of the dictionary object cache. Both caches store table definitions but serve different parts of the MySQL server. Objects in one cache have no dependence on the existence objects in the other. For more information, see Section 14.4, “Dictionary Object Cache”.

- `table_open_cache`

System Variable	Name	table_open_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	2000
	Min Value	1
	Max Value	524288

The number of open tables for all threads. Increasing this value increases the number of file descriptors that `mysqld` requires. You can check whether you need to increase the table cache by checking the `Opened_tables` status variable. See Section 5.1.7, “Server Status Variables”. If the value of `Opened_tables` is large and you do not use `FLUSH TABLES` often (which just forces all tables to be closed and reopened), then you should increase the value of the `table_open_cache` variable. For more information about the table cache, see Section 8.4.3.1, “How MySQL Opens and Closes Tables”.

- `table_open_cache_instances`

System Variable	Name	table_open_cache_instances
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	16
	Min Value	1
	Max Value	64

The number of open tables cache instances. To improve scalability by reducing contention among sessions, the open tables cache can be partitioned into several smaller cache instances of size `table_open_cache` / `table_open_cache_instances` . A session needs to lock only one instance to access it for DML statements. This segments cache access among instances, permitting higher performance for operations that use the cache when there are many sessions accessing tables. (DDL statements still require a lock on the entire cache, but such statements are much less frequent than DML statements.)

A value of 8 or 16 is recommended on systems that routinely use 16 or more cores.

- `temptable_max_ram`

Introduced	8.0.2	
Command-Line Format	`--temptable-max-ram=#`	
System Variable	**Name**	`temptable_max_ram`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1073741824`
	Min Value	`2097152`
	Max Value	`2^64-1`

Defines the maximum amount of memory that the TempTable storage engine can use before data is stored to disk. The default value is 1073741824 bits (1GiB). For more information, see Section 8.4.4, “Internal Temporary Table Use in MySQL”.

- `thread_cache_size`

Command-Line Format	`--thread-cache-size=#`	
System Variable	**Name**	`thread_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`
	Min Value	`0`
	Max Value	`16384`

How many threads the server should cache for reuse. When a client disconnects, the client's threads are put in the cache if there are fewer than `thread_cache_size` threads there. Requests for threads are satisfied by reusing threads taken from the cache if possible, and only when the cache is empty is a new thread created. This variable can be increased to improve performance if you have a lot of new connections. Normally, this does not provide a notable performance improvement if you have a good thread implementation. However, if your server sees hundreds of connections per second you should normally set `thread_cache_size` high enough so that most

new connections use cached threads. By examining the difference between the `Connections` and `Threads_created` status variables, you can see how efficient the thread cache is. For details, see Section 5.1.7, “Server Status Variables”.

The default value is based on the following formula, capped to a limit of 100:

```
8 + (max_connections / 100)
```

- `thread_handling`

Command-Line Format	`--thread-handling=name`	
System Variable	**Name**	`thread_handling`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	enumeration
	Default	`one-thread-per-connection`
	Valid Values	`no-threads`
		`one-thread-per-connection`
		`dynamically-loaded`

The thread-handling model used by the server for connection threads. The permissible values are `no-threads` (the server uses a single thread to handle one connection) and `one-thread-per-connection` (the server uses one thread to handle each client connection). `no-threads` is useful for debugging under Linux; see Section 28.5, “Debugging and Porting MySQL”.

- `thread_stack`

Command-Line Format	`--thread-stack=#`	
System Variable	**Name**	`thread_stack`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`196608`
	Min Value	`131072`
	Max Value	`4294967295`
	Block Size	`1024`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`262144`
	Min Value	`131072`
	Max Value	`18446744073709551615`

	Block Size	1024

The stack size for each thread. The default of 192KB (256KB for 64-bit systems) is large enough for normal operation. If the thread stack size is too small, it limits the complexity of the SQL statements that the server can handle, the recursion depth of stored procedures, and other memory-consuming actions.

- `time_format`

  This variable is unused. It is deprecated and will be removed in a future MySQL release.

- `time_zone`

System Variable	**Name**	`time_zone`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The current time zone. This variable is used to initialize the time zone for each client that connects. By default, the initial value of this is `'SYSTEM'` (which means, "use the value of `system_time_zone`"). The value can be specified explicitly at server startup with the `--default-time-zone` option. See Section 10.6, "MySQL Server Time Zone Support".

- `timestamp`

System Variable	**Name**	`timestamp`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	**Type**	numeric

Set the time for this client. This is used to get the original timestamp if you use the binary log to restore rows. *`timestamp_value`* should be a Unix epoch timestamp (a value like that returned by `UNIX_TIMESTAMP()`, not a value in `'YYYY-MM-DD hh:mm:ss'` format) or `DEFAULT`.

Setting `timestamp` to a constant value causes it to retain that value until it is changed again. Setting `timestamp` to `DEFAULT` causes its value to be the current date and time as of the time it is accessed.

In MySQL 8.0, `timestamp` is a `DOUBLE` rather than `BIGINT` because its value includes a microseconds part.

`SET timestamp` affects the value returned by `NOW()` but not by `SYSDATE()`. This means that timestamp settings in the binary log have no effect on invocations of `SYSDATE()`. The server can be started with the `--sysdate-is-now` option to cause `SYSDATE()` to be an alias for `NOW()`, in which case `SET timestamp` affects both functions.

- `tls_version`

	Command-Line Format	`--tls-version=protocol_list`
System Variable	**Name**	`tls_version`

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (OpenSSL)	**Type**	string
	Default	`TLSv1,TLSv1.1,TLSv1.2`
Permitted Values (yaSSL)	**Type**	string
	Default	`TLSv1,TLSv1.1`

The protocols permitted by the server for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this variable depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, "Secure Connection Protocols and Ciphers".

- `tmp_table_size`

Command-Line Format	`--tmp-table-size=#`	
System Variable	**Name**	`tmp_table_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`16777216`
	Min Value	`1024`
	Max Value	`18446744073709551615`

The maximum size of internal in-memory temporary tables. This variable does not apply to user-created `MEMORY` tables.

The actual limit is determined from whichever of the values of `tmp_table_size` and `max_heap_table_size` is smaller. If an in-memory temporary table exceeds the limit, MySQL automatically converts it to an on-disk temporary table. The `internal_tmp_disk_storage_engine` option defines the storage engine used for on-disk temporary tables.

Increase the value of `tmp_table_size` (and `max_heap_table_size` if necessary) if you do many advanced `GROUP BY` queries and you have lots of memory.

You can compare the number of internal on-disk temporary tables created to the total number of internal temporary tables created by comparing the values of the `Created_tmp_disk_tables` and `Created_tmp_tables` variables.

See also Section 8.4.4, "Internal Temporary Table Use in MySQL".

- `tmpdir`

Command-Line Format	`--tmpdir=dir_name`	
System Variable	**Name**	`tmpdir`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	directory name

The directory used for temporary files and temporary tables. This variable can be set to a list of several paths that are used in round-robin fashion. Paths should be separated by colon characters (`:`) on Unix and semicolon characters (`;`) on Windows.

The multiple-directory feature can be used to spread the load between several physical disks. If the MySQL server is acting as a replication slave, you should not set `tmpdir` to point to a directory on a memory-based file system or to a directory that is cleared when the server host restarts. A replication slave needs some of its temporary files to survive a machine restart so that it can replicate temporary tables or `LOAD DATA INFILE` operations. If files in the temporary file directory are lost when the server restarts, replication fails. You can set the slave's temporary directory using the `slave_load_tmpdir` variable. In that case, the slave will not use the general `tmpdir` value and you can set `tmpdir` to a nonpermanent location.

- `transaction_alloc_block_size`

Command-Line Format	`--transaction-alloc-block-size=#`	
System Variable	**Name**	`transaction_alloc_block_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`8192`
	Min Value	`1024`
	Max Value	`131072`
	Block Size	`1024`

The amount in bytes by which to increase a per-transaction memory pool which needs memory. See the description of `transaction_prealloc_size`.

- `transaction_prealloc_size`

Command-Line Format	`--transaction-prealloc-size=#`	
System Variable	**Name**	`transaction_prealloc_size`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`4096`
	Min Value	`1024`
	Max Value	`131072`

	Block Size	1024

There is a per-transaction memory pool from which various transaction-related allocations take memory. The initial size of the pool in bytes is `transaction_prealloc_size`. For every allocation that cannot be satisfied from the pool because it has insufficient memory available, the pool is increased by `transaction_alloc_block_size` bytes. When the transaction ends, the pool is truncated to `transaction_prealloc_size` bytes.

By making `transaction_prealloc_size` sufficiently large to contain all statements within a single transaction, you can avoid many `malloc()` calls.

- `transaction_write_set_extraction`

Command-Line Format	`--transaction-write-set-extraction=[value]`	
System Variable	**Name**	`transaction_write_set_extraction`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`OFF`
	Valid Values	`OFF`
		`MURMUR32`
		`XXHASH64`
Permitted Values (>= 8.0.2)	**Type**	enumeration
	Default	`XXHASH64`
	Valid Values	`OFF`
		`MURMUR32`
		`XXHASH64`

Defines the algorithm used to hash the writes extracted during a transaction. If you are using Group Replication, this variable must be set to `XXHASH64` because the process of extracting the writes from a transaction is required for conflict detection on all group members. See Section 21.7.1, “Group Replication Requirements”.

Note

The value of this variable cannot be changed when `binlog_transaction_dependency_tracking` is set to either of `WRITESET` or `WRITESET_SESSION`.

- `tx_isolation`

System Variable	**Name**	`tx_isolation`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`REPEATABLE-READ`

	Valid Values	`READ-UNCOMMITTED`
		`READ-COMMITTED`
		`REPEATABLE-READ`
		`SERIALIZABLE`

The default transaction isolation level. Defaults to `REPEATABLE-READ`.

This variable can be set directly, or indirectly using the `SET TRANSACTION` statement. See Section 13.3.6, “SET TRANSACTION Syntax”. If you set `tx_isolation` directly to an isolation level name that contains a space, the name should be enclosed within quotation marks, with the space replaced by a dash. For example:

```
SET tx_isolation = 'READ-COMMITTED';
```

Any unique prefix of a valid value may be used to set the value of this variable.

The default transaction isolation level can also be set at startup using the `--transaction-isolation` server option.

- `tx_read_only`

System Variable	**Name**	`tx_read_only`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

The default transaction access mode. The value can be `OFF` (read/write, the default) or `ON` (read only).

This variable can be set directly, or indirectly using the `SET TRANSACTION` statement. See Section 13.3.6, “SET TRANSACTION Syntax”.

To set the default transaction access mode at startup, use the `--transaction-read-only` server option.

- `unique_checks`

System Variable	**Name**	`unique_checks`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`1`

If set to 1 (the default), uniqueness checks for secondary indexes in `InnoDB` tables are performed. If set to 0, storage engines are permitted to assume that duplicate keys are not present in input data. If you know for certain that your data does not contain uniqueness violations, you can set this to 0 to speed up large table imports to `InnoDB`.

Setting this variable to 0 does not *require* storage engines to ignore duplicate keys. An engine is still permitted to check for them and issue duplicate-key errors if it detects them.

- `updatable_views_with_limit`

Command-Line Format	`--updatable-views-with-limit=#`	
System Variable	Name	`updatable_views_with_limit`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`1`

This variable controls whether updates to a view can be made when the view does not contain all columns of the primary key defined in the underlying table, if the update statement contains a `LIMIT` clause. (Such updates often are generated by GUI tools.) An update is an `UPDATE` or `DELETE` statement. Primary key here means a `PRIMARY KEY`, or a `UNIQUE` index in which no column can contain `NULL`.

The variable can have two values:

- `1` or `YES`: Issue a warning only (not an error message). This is the default value.
- `0` or `NO`: Prohibit the update.

- `validate_password_xxx`

The `validate_password` plugin implements a set of system variables having names of the form `validate_password_xxx`. These variables affect password testing by that plugin; see Section 6.5.3.2, “Password Validation Plugin Options and Variables”.

- `validate_user_plugins`

System Variable	Name	`validate_user_plugins`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`ON`

If this variable is enabled (the default), the server checks each user account and produces a warning if conditions are found that would make the account unusable:

- The account requires an authentication plugin that is not loaded.
- The account requires the `sha256_password` authentication plugin but the server was started with neither SSL nor RSA enabled as required by this plugin.

Enabling `validate_user_plugins` slows down server initialization and `FLUSH PRIVILEGES`. If you do not require the additional checking, you can disable this variable at startup to avoid the performance decrement.

- `version`

The version number for the server. The value might also include a suffix indicating server build or configuration information. `-log` indicates that one or more of the general log, slow query log, or binary log are enabled. `-debug` indicates that the server was built with debugging support enabled.

- `version_comment`

System Variable	Name	version_comment
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The `CMake` configuration program has a `COMPILATION_COMMENT` option that permits a comment to be specified when building MySQL. This variable contains the value of that comment. See Section 2.8.4, “MySQL Source-Configuration Options”.

- `version_compile_machine`

System Variable	Name	version_compile_machine
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The type of the server binary.

- `version_compile_os`

System Variable	Name	version_compile_os
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The type of operating system on which MySQL was built.

- `wait_timeout`

Command-Line Format	--wait-timeout=#	
System Variable	Name	wait_timeout
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values (Windows)	Type	integer
	Default	28800
	Min Value	1

	Max Value	2147483
Permitted Values (Other)	**Type**	integer
	Default	28800
	Min Value	1
	Max Value	31536000

The number of seconds the server waits for activity on a noninteractive connection before closing it.

On thread startup, the session `wait_timeout` value is initialized from the global `wait_timeout` value or from the global `interactive_timeout` value, depending on the type of client (as defined by the `CLIENT_INTERACTIVE` connect option to `mysql_real_connect()`). See also `interactive_timeout`.

- `warning_count`

  The number of errors, warnings, and notes that resulted from the last statement that generated messages. This variable is read only. See Section 13.7.5.40, “SHOW WARNINGS Syntax”.

5.1.6 Using System Variables

The MySQL server maintains many system variables that indicate how it is configured. Section 5.1.5, “Server System Variables”, describes the meaning of these variables. Each system variable has a default value. System variables can be set at server startup using options on the command line or in an option file. Most of them can be changed dynamically while the server is running by means of the `SET` statement, which enables you to modify operation of the server without having to stop and restart it. You can also use system variable values in expressions.

There are two scopes in which system variables exist. Global variables affect the overall operation of the server. Session variables affect its operation for individual client connections. A given system variable can have both a global and a session value. Global and session system variables are related as follows:

- When the server starts, it initializes each global variable to its default value. These defaults can be changed by options specified on the command line or in an option file. (See Section 4.2.3, “Specifying Program Options”.)
- The server also maintains a set of session variables for each client that connects. The client's session variables are initialized at connect time using the current values of the corresponding global variables. For example, a client's SQL mode is controlled by the session `sql_mode` value, which is initialized when the client connects to the value of the global `sql_mode` value.

  For some system variables, the session value is not initialized from the corresponding global value; if so, that is indicated in the variable description.

System variable values can be set globally at server startup by using options on the command line or in an option file. When you use a startup option to set a variable that takes a numeric value, the value can be given with a suffix of `K`, `M`, or `G` (either uppercase or lowercase) to indicate a multiplier of 1024, 1024^2 or 1024^3; that is, units of kilobytes, megabytes, or gigabytes, respectively. Thus, the following command starts the server with a query cache size of 16 megabytes and a maximum packet size of one gigabyte:

```
mysqld --query_cache_size=16M --max_allowed_packet=1G
```

Within an option file, those variables are set like this:

```
```

```
[mysqld]
query_cache_size=16M
max_allowed_packet=1G
```

The lettercase of suffix letters does not matter; `16M` and `16m` are equivalent, as are `1G` and `1g`.

To restrict the maximum value to which a system variable can be set at runtime with the `SET` statement, specify this maximum by using an option of the form `--maximum-var_name=value` at server startup. For example, to prevent the value of `query_cache_size` from being increased to more than 32MB at runtime, use the option `--maximum-query_cache_size=32M`.

Many system variables are dynamic and can be changed at runtime by using the `SET` statement. For a list, see Section 5.1.6.2, “Dynamic System Variables”. To change a system variable with `SET`, refer to it by name, optionally preceded by a modifier:

- To indicate that a variable is a global variable, precede its name by the `GLOBAL` keyword or the `@@global.` qualifier:

  ```
  SET GLOBAL max_connections = 1000;
  SET @@global.max_connections = 1000;
  ```

  The `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege is required to set global variables.

- Another way to set a global variable is to precede its name by the `PERSIST` keyword or the `@@persist.` qualifier:

  ```
  SET PERSIST max_connections = 1000;
  SET @@persist.max_connections = 1000;
  ```

  This `SET` syntax enables you to make configuration changes at runtime that also persist across server restarts. Like `SET GLOBAL`, `SET PERSIST` changes the runtime variable value, but also writes the variable setting to an option file named `mysqld-auto.cnf` in the data directory (replacing any existing variable setting if there is one). At startup, the server processes this file after all other option files. The `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege is required to persist global variables.

  Note

  Management of the `mysqld-auto.cnf` file should be left to the server and not performed manually:

  - Removal of the file results in a loss of all persisted settings at the next server startup. (This is permissible if your intent is to reconfigure the server without these settings.) To remove all settings in the file without removing the file itself, use this statement:

    ```
    RESET PERSIST;
    ```

  - Manual changes to the file may result in a parse error at server startup. In this case, the server reports an error and exits. If this issue occurs, start the server with the `persisted_globals_load` system variable disabled or with the `--no-defaults` option. Alternatively, remove the `mysqld-auto.cnf` file, but, as noted previously, removing this file results in a loss of all persisted settings.

  A plugin variable can be persisted if the plugin is installed when `SET PERSIST` is executed. Assignment of the persisted plugin variable takes effect for subsequent server restarts if the plugin is still installed. If the plugin is no longer installed, the plugin variable will not exist when the server reads the `mysqld-auto.cnf` file. In this case, the server writes a warning to the error log and continues:

```
currently unknown variable 'var_name'
was read from the persisted config file
```

- To indicate that a variable is a session variable, precede its name by the `SESSION` keyword or either the `@@session.` or `@@` qualifier:

```
SET SESSION sql_mode = 'TRADITIONAL';
SET @@session.sql_mode = 'TRADITIONAL';
SET @@sql_mode = 'TRADITIONAL';
```

  Setting a session variable normally requires no special privilege, although there are exceptions that require the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege (such as `sql_log_bin`). A client can change its own session variables, but not those of any other client.

  Session-only system variables cannot be persisted. They cannot be set at server startup, so there is no reason to list them in `mysqld-auto.cnf`.

- `LOCAL` and `@@local.` are synonyms for `SESSION` and `@@session.`.
- If no modifier is present, `SET` changes the session variable.
- An error occurs under these circumstances:
  - Use of `SET GLOBAL` (or `@@global.`) or `SET PERSIST` (or `@@persist.`) when setting a variable that has only a session value:

```
mysql> SET GLOBAL sql_log_bin = ON;
ERROR 1231 (42000): Variable 'sql_log_bin' can't be
set to the value of 'ON'
```

  - Omission of `GLOBAL` (or `@@global.`) or `PERSIST` (or `@@persist.`) when setting a variable that has only a global value:

```
mysql> SET max_connections = 1000;
ERROR 1229 (HY000): Variable 'max_connections' is a
GLOBAL variable and should be set with SET GLOBAL
```

  - Use of `SET SESSION` (or `@@SESSION.`) when setting a variable that has only a global value:

```
mysql> SET SESSION max_connections = 1000;
ERROR 1229 (HY000): Variable 'max_connections' is a
GLOBAL variable and should be set with SET GLOBAL
```

The preceding modifiers apply only to system variables. An error occurs for attempts to apply them to user-defined variables, stored procedure or function parameters, or stored program local variables.

A `SET` statement can contain multiple variable assignments, separated by commas. This statement assigns values to a user-defined variable and a system variable:

```
SET @x = 1, SESSION sql_mode = '';
```

If you set multiple system variables, the most recent `GLOBAL`, `PERSIST`, or `SESSION` modifier in the statement is used for following assignments that have no modifier specified.

Examples of multiple-variable assignment:

```
SET GLOBAL sort_buffer_size = 1000000, SESSION sort_buffer_size = 1000000;
SET @@global.sort_buffer_size = 1000000, @@local.sort_buffer_size = 1000000;
SET GLOBAL max_connections = 1000, sort_buffer_size = 1000000;
```

If any variable assignment in a SET statement fails, the entire statement fails and no variables are changed, nor is the mysqld-auto.cnf file changed.

If you change a session system variable, the value remains in effect within your session until you change the variable to a different value or the session ends. The change has no effect on other sessions.

If you change a global system variable, the value is remembered and used for new sessions until you change the variable to a different value or the server exits. The change is visible to any client that accesses the global variable. However, the change affects the corresponding session variable only for clients that connect after the change. The global variable change does not affect the session variable for any current client sessions (not even the session within which the SET GLOBAL statement occurred).

To make a global system variable setting permanent so that it applies across server restarts, modify it with SET PERSIST to record it in the mysqld-auto.cnf file. It is also possible to use SET GLOBAL and manually modify a my.cnf file, but that is more cumbersome, and an error in a manually entered setting might not be discovered until much later. SET PERSIST is more convenient and avoids the possibility of malformed settings.

The Performance Schema persisted_variables table provides an SQL interface to the mysqld-auto.cnf file, enabling its contents to be inspected at runtime using SELECT statements. See Section 25.11.13.1, “Performance Schema persisted_variables Table”.

The Performance Schema variables_info table contains information showing when and by which user each system variable was most recently set. See Section 25.11.13.2, “Performance Schema variables_info Table”.

To set a GLOBAL value to the compiled-in MySQL default value or a SESSION variable to the current corresponding GLOBAL value, set the variable to the value DEFAULT. For example, the following two statements are identical in setting the session value of max_join_size to the current global value:

```
SET @@session.max_join_size=DEFAULT;
SET @@session.max_join_size=@@global.max_join_size;
```

Not all system variables can be set to DEFAULT. In such cases, assigning DEFAULT results in an error.

With SET PERSIST (or @@persist.), setting a global variable to DEFAULT or to its literal default value assigns the variable its default value and adds a setting for the variable to mysqld-auto.cnf. To remove the variable from the file, use RESET PERSIST.

An error occurs for attempts to assign DEFAULT to user-defined variables, stored procedure or function parameters, or stored program local variables.

To refer to the value of a system variable in expressions, use one of the @@-modifiers (except @@persist., which is not permitted in expressions). For example, you can retrieve values in a SELECT statement like this:

```
SELECT @@global.sql_mode, @@session.sql_mode, @@sql_mode;
```

For a reference to a system variable in an expression as *@@var_name* (rather than with @@global. or @@session.), MySQL returns the session value if it exists and the global value otherwise. This differs from SET *@@var_name = expr*, which always refers to the session value.

Suffixes for specifying a value multiplier can be used when setting a variable at server startup, but not to set the value with SET at runtime. On the other hand, with SET you can assign a variable's value using an expression, which is not true when you set a variable at server startup. For example, the first of the following lines is legal at server startup, but the second is not:

```
shell> mysql --max_allowed_packet=16M
shell> mysql --max_allowed_packet=16*1024*1024
```

Conversely, the second of the following lines is legal at runtime, but the first is not:

```
mysql> SET GLOBAL max_allowed_packet=16M;
mysql> SET GLOBAL max_allowed_packet=16*1024*1024;
```

Note

Some system variables can be enabled with the `SET` statement by setting them to `ON` or `1`, or disabled by setting them to `OFF` or `0`. However, to set such a variable on the command line or in an option file, you must set it to `1` or `0`; setting it to `ON` or `OFF` will not work. For example, on the command line, `--delay_key_write=1` works but `--delay_key_write=ON` does not.

To display system variable names and values, use the `SHOW VARIABLES` statement:

```
mysql> SHOW VARIABLES;
+---------------------------------+-----------------------------------+
| Variable_name                   | Value                             |
+---------------------------------+-----------------------------------+
| auto_increment_increment        | 1                                 |
| auto_increment_offset           | 1                                 |
| automatic_sp_privileges         | ON                                |
| back_log                        | 50                                |
| basedir                         | /home/mysql/                      |
| binlog_cache_size               | 32768                             |
| bulk_insert_buffer_size         | 8388608                           |
| character_set_client            | utf8                              |
| character_set_connection        | utf8                              |
| character_set_database          | utf8mb4                           |
| character_set_filesystem        | binary                            |
| character_set_results           | utf8                              |
| character_set_server            | utf8mb4                           |
| character_set_system            | utf8                              |
| character_sets_dir              | /home/mysql/share/mysql/charsets/ |
| collation_connection            | utf8_general_ci                   |
| collation_database              | utf8mb4_0900_ai_ci                |
| collation_server                | utf8mb4_0900_ai_ci                |
...
| innodb_autoextend_increment     | 8                                 |
| innodb_buffer_pool_size         | 8388608                           |
| innodb_commit_concurrency       | 0                                 |
| innodb_concurrency_tickets      | 500                               |
| innodb_data_file_path           | ibdata1:10M:autoextend            |
| innodb_data_home_dir            |                                   |
...
| version                         | 8.0.1-dmr-log                     |
| version_comment                 | Source distribution               |
| version_compile_machine         | i686                              |
| version_compile_os              | suse-linux                        |
| wait_timeout                    | 28800                             |
+---------------------------------+-----------------------------------+
```

With a `LIKE` clause, the statement displays only those variables that match the pattern. To obtain a specific variable name, use a `LIKE` clause as shown:

```
SHOW VARIABLES LIKE 'max_join_size';
SHOW SESSION VARIABLES LIKE 'max_join_size';
```

To get a list of variables whose name match a pattern, use the `%` wildcard character in a `LIKE` clause:

```
SHOW VARIABLES LIKE '%size%';
SHOW GLOBAL VARIABLES LIKE '%size%';
```

Wildcard characters can be used in any position within the pattern to be matched. Strictly speaking, because `_` is a wildcard that matches any single character, you should escape it as `\_` to match it literally. In practice, this is rarely necessary.

For `SHOW VARIABLES`, if you specify neither `GLOBAL` nor `SESSION`, MySQL returns `SESSION` values.

The reason for requiring the `GLOBAL` keyword when setting `GLOBAL`-only variables but not when retrieving them is to prevent problems in the future:

- Were we to remove a `SESSION` variable that has the same name as a `GLOBAL` variable, a client with the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege might accidentally change the `GLOBAL` variable rather than just the `SESSION` variable for its own connection.
- Were we to add a `SESSION` variable with the same name as a `GLOBAL` variable, a client that intends to change the `GLOBAL` variable might find only its own `SESSION` variable changed.

5.1.6.1 Structured System Variables

A structured variable differs from a regular system variable in two respects:

- Its value is a structure with components that specify server parameters considered to be closely related.
- There might be several instances of a given type of structured variable. Each one has a different name and refers to a different resource maintained by the server.

MySQL supports one structured variable type, which specifies parameters governing the operation of key caches. A key cache structured variable has these components:

- `key_buffer_size`
- `key_cache_block_size`
- `key_cache_division_limit`
- `key_cache_age_threshold`

This section describes the syntax for referring to structured variables. Key cache variables are used for syntax examples, but specific details about how key caches operate are found elsewhere, in Section 8.10.2, “The MyISAM Key Cache”.

To refer to a component of a structured variable instance, you can use a compound name in *`instance_name.component_name`* format. Examples:

```
hot_cache.key_buffer_size
hot_cache.key_cache_block_size
cold_cache.key_cache_block_size
```

For each structured system variable, an instance with the name of `default` is always predefined. If you refer to a component of a structured variable without any instance name, the `default` instance is used. Thus, `default.key_buffer_size` and `key_buffer_size` both refer to the same system variable.

Structured variable instances and components follow these naming rules:

- For a given type of structured variable, each instance must have a name that is unique *within* variables of that type. However, instance names need not be unique *across* structured variable types. For example, each structured variable has an instance named `default`, so `default` is not unique across variable types.
- The names of the components of each structured variable type must be unique across all system variable names. If this were not true (that is, if two different types of structured variables could share component member names), it would not be clear which default structured variable to use for references to member names that are not qualified by an instance name.
- If a structured variable instance name is not legal as an unquoted identifier, refer to it as a quoted identifier using backticks. For example, `hot-cache` is not legal, but `` `hot-cache` `` is.

- `global`, `session`, and `local` are not legal instance names. This avoids a conflict with notation such as `@@global.var_name` for referring to nonstructured system variables.

Currently, the first two rules have no possibility of being violated because the only structured variable type is the one for key caches. These rules will assume greater significance if some other type of structured variable is created in the future.

With one exception, you can refer to structured variable components using compound names in any context where simple variable names can occur. For example, you can assign a value to a structured variable using a command-line option:

```
shell> mysqld --hot_cache.key_buffer_size=64K
```

In an option file, use this syntax:

```
[mysqld]
hot_cache.key_buffer_size=64K
```

If you start the server with this option, it creates a key cache named `hot_cache` with a size of 64KB in addition to the default key cache that has a default size of 8MB.

Suppose that you start the server as follows:

```
shell> mysqld --key_buffer_size=256K \
         --extra_cache.key_buffer_size=128K \
         --extra_cache.key_cache_block_size=2048
```

In this case, the server sets the size of the default key cache to 256KB. (You could also have written `--default.key_buffer_size=256K`.) In addition, the server creates a second key cache named `extra_cache` that has a size of 128KB, with the size of block buffers for caching table index blocks set to 2048 bytes.

The following example starts the server with three different key caches having sizes in a 3:1:1 ratio:

```
shell> mysqld --key_buffer_size=6M \
         --hot_cache.key_buffer_size=2M \
         --cold_cache.key_buffer_size=2M
```

Structured variable values may be set and retrieved at runtime as well. For example, to set a key cache named `hot_cache` to a size of 10MB, use either of these statements:

```
mysql> SET GLOBAL hot_cache.key_buffer_size = 10*1024*1024;
mysql> SET @@global.hot_cache.key_buffer_size = 10*1024*1024;
```

To retrieve the cache size, do this:

```
mysql> SELECT @@global.hot_cache.key_buffer_size;
```

However, the following statement does not work. The variable is not interpreted as a compound name, but as a simple string for a `LIKE` pattern-matching operation:

```
mysql> SHOW GLOBAL VARIABLES LIKE 'hot_cache.key_buffer_size';
```

This is the exception to being able to use structured variable names anywhere a simple variable name may occur.

5.1.6.2 Dynamic System Variables

Many server system variables are dynamic and can be set at runtime using `SET GLOBAL` or `SET SESSION`. You can also obtain their values using `SELECT`. See Section 5.1.6, “Using System Variables”.

The following table shows the full list of all dynamic system variables. The last column indicates for each variable whether GLOBAL or SESSION (or both) apply. The table also lists session options that can be set with the SET statement. Section 5.1.5, “Server System Variables”, discusses these options.

Variables that have a type of “string” take a string value. Variables that have a type of “numeric” take a numeric value. Variables that have a type of “boolean” can be set to 0, 1, ON or OFF. Variables that are marked as “enumeration” normally should be set to one of the available values for the variable, but can also be set to the number that corresponds to the desired enumeration value. For enumerated system variables, the first enumeration value corresponds to 0. This differs from ENUM columns, for which the first enumeration value corresponds to 1.

Table 5.3 Dynamic Variable Summary

Variable Name	Variable Type	Variable Scope
auto_increment_increment	integer	GLOBAL | SESSION
auto_increment_offset	integer	GLOBAL | SESSION
autocommit	boolean	GLOBAL | SESSION
automatic_sp_privileges	boolean	GLOBAL
avoid_temporal_upgrade	boolean	GLOBAL
big_tables	boolean	GLOBAL | SESSION
binlog_cache_size	integer	GLOBAL
binlog_checksum	string	GLOBAL
binlog_direct_non_transactional_updates	boolean	GLOBAL | SESSION
binlog_error_action	enumeration	GLOBAL | SESSION
binlog_expire_logs_seconds	integer	GLOBAL
binlog_format	enumeration	GLOBAL | SESSION
binlog_group_commit_sync_delay	integer	GLOBAL
binlog_group_commit_sync_no_delay_count	integer	GLOBAL
binlog_max_flush_queue_time	integer	GLOBAL
binlog_order_commits	boolean	GLOBAL
binlog_row_image=image_type	enumeration	GLOBAL | SESSION
binlog_row_metadata=metadata_type	enumeration	GLOBAL
binlog_rows_query_log_events	boolean	GLOBAL | SESSION
binlog_stmt_cache_size	integer	GLOBAL
binlog_transaction_dependency_history_size	integer	GLOBAL
binlog_transaction_dependency_tracking	enumeration	GLOBAL
binlogging_impossible_mode	enumeration	GLOBAL | SESSION
block_encryption_mode	string	GLOBAL | SESSION
bulk_insert_buffer_size	integer	GLOBAL | SESSION
character_set_client	string	GLOBAL | SESSION
character_set_connection	string	GLOBAL | SESSION
character_set_database	string	GLOBAL | SESSION
character_set_filesystem	string	GLOBAL | SESSION
character_set_results	string	GLOBAL | SESSION
character_set_server	string	GLOBAL | SESSION
check_proxy_users	boolean	GLOBAL

Variable Name	Variable Type	Variable Scope
collation_connection	string	GLOBAL | SESSION
collation_database	string	GLOBAL | SESSION
collation_server	string	GLOBAL | SESSION
completion_type	enumeration	GLOBAL | SESSION
concurrent_insert	enumeration	GLOBAL
connect_timeout	integer	GLOBAL
connection_control_failed_connections_threshold	integer	GLOBAL
connection_control_max_connection_delay	integer	GLOBAL
connection_control_min_connection_delay	integer	GLOBAL
debug	string	GLOBAL | SESSION
debug_sync	string	SESSION
default_password_lifetime	integer	GLOBAL
default_storage_engine	enumeration	GLOBAL | SESSION
default_tmp_storage_engine	enumeration	GLOBAL | SESSION
default_week_format	integer	GLOBAL | SESSION
delay_key_write	enumeration	GLOBAL
delayed_insert_limit	integer	GLOBAL
delayed_insert_timeout	integer	GLOBAL
delayed_queue_size	integer	GLOBAL
div_precision_increment	integer	GLOBAL | SESSION
end_markers_in_json	boolean	GLOBAL | SESSION
enforce_gtid_consistency	enumeration	GLOBAL
enforce_gtid_consistency	enumeration	GLOBAL
eq_range_index_dive_limit	integer	GLOBAL | SESSION
event_scheduler	enumeration	GLOBAL
executed_gtids_compression_period	integer	GLOBAL
expire_logs_days	integer	GLOBAL
explicit_defaults_for_timestamp	boolean	GLOBAL | SESSION
flush	boolean	GLOBAL
flush_time	integer	GLOBAL
foreign_key_checks	boolean	GLOBAL | SESSION
ft_boolean_syntax	string	GLOBAL
general_log	boolean	GLOBAL
general_log_file	filename	GLOBAL
group_concat_max_len	integer	GLOBAL | SESSION
group_replication_allow_local_disjoint_gtids_join	boolean	GLOBAL
group_replication_allow_local_lower_version_join	boolean	GLOBAL
group_replication_auto_increment_increment	integer	GLOBAL
group_replication_bootstrap_group	boolean	GLOBAL
group_replication_components_stop_timeout	integer	GLOBAL
group_replication_compression_threshold	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
group_replication_enforce_update_everywhere_checks	boolean	GLOBAL
group_replication_flow_control_applier_threshold	integer	GLOBAL
group_replication_flow_control_certifier_threshold	integer	GLOBAL
group_replication_flow_control_hold_percent	integer	GLOBAL
group_replication_flow_control_max_commit_quota	integer	GLOBAL
group_replication_flow_control_member_quota_percent	integer	GLOBAL
group_replication_flow_control_min_quota	integer	GLOBAL
group_replication_flow_control_min_recovery_quota	integer	GLOBAL
group_replication_flow_control_mode	enumeration	GLOBAL
group_replication_flow_control_period	integer	GLOBAL
group_replication_flow_control_release_percent	integer	GLOBAL
group_replication_force_members	string	GLOBAL
group_replication_group_name	string	GLOBAL
group_replication_group_seeds	string	GLOBAL
group_replication_gtid_assignment_block_size	integer	GLOBAL
group_replication_ip_whitelist	string	GLOBAL
group_replication_local_address	string	GLOBAL
group_replication_member_weight	integer	GLOBAL
group_replication_poll_spin_loops	integer	GLOBAL
group_replication_recovery_complete_at	enumeration	GLOBAL
group_replication_recovery_reconnect_interval	integer	GLOBAL
group_replication_recovery_retry_count	integer	GLOBAL
group_replication_recovery_ssl_ca	string	GLOBAL
group_replication_recovery_ssl_capath	string	GLOBAL
group_replication_recovery_ssl_cert	string	GLOBAL
group_replication_recovery_ssl_cipher	string	GLOBAL
group_replication_recovery_ssl_crl	string	GLOBAL
group_replication_recovery_ssl_crlpath	string	GLOBAL
group_replication_recovery_ssl_key	string	GLOBAL
group_replication_recovery_ssl_verify_server_cert	boolean	GLOBAL
group_replication_recovery_use_ssl	boolean	GLOBAL
group_replication_single_primary_mode	boolean	GLOBAL
group_replication_ssl_mode	enumeration	GLOBAL
group_replication_start_on_boot	boolean	GLOBAL
group_replication_transaction_size_limit	integer	GLOBAL
group_replication_unreachable_majority_timeout	integer	GLOBAL
gtid_executed_compression_period	integer	GLOBAL
gtid_mode	enumeration	GLOBAL
gtid_mode	enumeration	GLOBAL
gtid_next	enumeration	SESSION
gtid_purged	string	GLOBAL

Variable Name	Variable Type	Variable Scope
histogram_generation_max_mem_size	integer	GLOBAL | SESSION
host_cache_size	integer	GLOBAL
identity	integer	SESSION
information_schema_stats	enumeration	GLOBAL | SESSION
init_connect	string	GLOBAL
init_slave	string	GLOBAL
innodb_adaptive_flushing	boolean	GLOBAL
innodb_adaptive_flushing_lwm	integer	GLOBAL
innodb_adaptive_hash_index	boolean	GLOBAL
innodb_adaptive_max_sleep_delay	integer	GLOBAL
innodb_api_bk_commit_interval	integer	GLOBAL
innodb_api_trx_level	integer	GLOBAL
innodb_autoextend_increment	integer	GLOBAL
innodb_background_drop_list_empty	boolean	GLOBAL
innodb_buffer_pool_dump_at_shutdown	boolean	GLOBAL
innodb_buffer_pool_dump_now	boolean	GLOBAL
innodb_buffer_pool_dump_pct	integer	GLOBAL
innodb_buffer_pool_filename	filename	GLOBAL
innodb_buffer_pool_load_abort	boolean	GLOBAL
innodb_buffer_pool_load_now	boolean	GLOBAL
innodb_buffer_pool_size	integer	GLOBAL
innodb_change_buffer_max_size	integer	GLOBAL
innodb_change_buffering	enumeration	GLOBAL
innodb_change_buffering_debug	integer	GLOBAL
innodb_checkpoint_disabled	boolean	GLOBAL
innodb_checksum_algorithm	enumeration	GLOBAL
innodb_cmp_per_index_enabled	boolean	GLOBAL
innodb_commit_concurrency	integer	GLOBAL
innodb_compress_debug	enumeration	GLOBAL
innodb_compression_failure_threshold_pct	integer	GLOBAL
innodb_compression_level	integer	GLOBAL
innodb_compression_pad_pct_max	integer	GLOBAL
innodb_concurrency_tickets	integer	GLOBAL
innodb_deadlock_detect	boolean	GLOBAL
innodb_default_row_format	enumeration	GLOBAL
innodb_disable_sort_file_cache	boolean	GLOBAL
innodb_fast_shutdown	integer	GLOBAL
innodb_fil_make_page_dirty_debug	integer	GLOBAL
innodb_file_per_table	boolean	GLOBAL
innodb_fill_factor	integer	GLOBAL
innodb_flush_log_at_timeout	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
innodb_flush_log_at_trx_commit	enumeration	GLOBAL
innodb_flush_neighbors	enumeration	GLOBAL
innodb_flush_sync	boolean	GLOBAL
innodb_flushing_avg_loops	integer	GLOBAL
innodb_ft_aux_table	string	GLOBAL
innodb_ft_enable_diag_print	boolean	GLOBAL
innodb_ft_enable_stopword	boolean	GLOBAL
innodb_ft_num_word_optimize	integer	GLOBAL
innodb_ft_result_cache_limit	integer	GLOBAL
innodb_ft_server_stopword_table	string	GLOBAL
innodb_ft_user_stopword_table	string	GLOBAL | SESSION
innodb_io_capacity	integer	GLOBAL
innodb_io_capacity_max	integer	GLOBAL
innodb_limit_optimistic_insert_debug	integer	GLOBAL
innodb_lock_wait_timeout	integer	GLOBAL | SESSION
innodb_log_checksums	boolean	GLOBAL
innodb_log_compressed_pages	boolean	GLOBAL
innodb_log_write_ahead_size	integer	GLOBAL
innodb_lru_scan_depth	integer	GLOBAL
innodb_max_dirty_pages_pct	numeric	GLOBAL
innodb_max_dirty_pages_pct_lwm	numeric	GLOBAL
innodb_max_purge_lag	integer	GLOBAL
innodb_max_purge_lag_delay	integer	GLOBAL
innodb_max_undo_log_size	integer	GLOBAL
innodb_merge_threshold_set_all_debug	integer	GLOBAL
innodb_monitor_disable	string	GLOBAL
innodb_monitor_enable	string	GLOBAL
innodb_monitor_reset	string	GLOBAL
innodb_monitor_reset_all	string	GLOBAL
innodb_old_blocks_pct	integer	GLOBAL
innodb_old_blocks_time	integer	GLOBAL
innodb_online_alter_log_max_size	integer	GLOBAL
innodb_optimize_fulltext_only	boolean	GLOBAL
innodb_print_all_deadlocks	boolean	GLOBAL
innodb_purge_batch_size	integer	GLOBAL
innodb_purge_rseg_truncate_frequency	integer	GLOBAL
innodb_random_read_ahead	boolean	GLOBAL
innodb_read_ahead_threshold	integer	GLOBAL
innodb_redo_log_encrypt	boolean	GLOBAL
innodb_replication_delay	integer	GLOBAL
innodb_rollback_segments	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
innodb_saved_page_number_debug	integer	GLOBAL
innodb_spin_wait_delay	integer	GLOBAL
innodb_stats_auto_recalc	boolean	GLOBAL
innodb_stats_include_delete_marked	boolean	GLOBAL
innodb_stats_method	enumeration	GLOBAL
innodb_stats_on_metadata	boolean	GLOBAL
innodb_stats_persistent	boolean	GLOBAL
innodb_stats_persistent_sample_pages	integer	GLOBAL
innodb_stats_transient_sample_pages	integer	GLOBAL
innodb_status_output	boolean	GLOBAL
innodb_status_output_locks	boolean	GLOBAL
innodb_strict_mode	boolean	GLOBAL | SESSION
innodb_sync_spin_loops	integer	GLOBAL
innodb_table_locks	boolean	GLOBAL | SESSION
innodb_thread_concurrency	integer	GLOBAL
innodb_thread_sleep_delay	integer	GLOBAL
innodb_tmpdir	dirname	GLOBAL | SESSION
innodb_trx_purge_view_update_only_debug	boolean	GLOBAL
innodb_trx_rseg_n_slots_debug	integer	GLOBAL
innodb_undo_log_encrypt	boolean	GLOBAL
innodb_undo_log_truncate	boolean	GLOBAL
innodb_undo_logs	integer	GLOBAL
innodb_undo_tablespaces	integer	GLOBAL
insert_id	integer	SESSION
interactive_timeout	integer	GLOBAL | SESSION
internal_tmp_disk_storage_engine	enumeration	GLOBAL
internal_tmp_mem_storage_engine	enumeration	SESSION
join_buffer_size	integer	GLOBAL | SESSION
keep_files_on_create	boolean	GLOBAL | SESSION
key_buffer_size	integer	GLOBAL
key_cache_age_threshold	integer	GLOBAL
key_cache_block_size	integer	GLOBAL
key_cache_division_limit	integer	GLOBAL
keyring_file_data	filename	GLOBAL
last_insert_id	integer	SESSION
lc_messages	string	GLOBAL | SESSION
lc_time_names	string	GLOBAL | SESSION
local_infile	boolean	GLOBAL
lock_wait_timeout	integer	GLOBAL | SESSION
log_bin_trust_function_creators	boolean	GLOBAL
log_builtin_as_identified_by_password	boolean	GLOBAL

Variable Name	Variable Type	Variable Scope
log_error_verbosity	integer	GLOBAL
log_output	set	GLOBAL
log_queries_not_using_indexes	boolean	GLOBAL
log_slow_admin_statements	boolean	GLOBAL
log_slow_slave_statements	boolean	GLOBAL
log_statements_unsafe_for_binlog	boolean	GLOBAL
log_syslog	boolean	GLOBAL
log_syslog_facility	string	GLOBAL
log_syslog_include_pid	boolean	GLOBAL
log_syslog_tag	string	GLOBAL
log_throttle_queries_not_using_indexes	integer	GLOBAL
log_timestamps	enumeration	GLOBAL
log_warnings	integer	GLOBAL
long_query_time	numeric	GLOBAL | SESSION
low_priority_updates	boolean	GLOBAL | SESSION
master_info_repository	string	GLOBAL
master_verify_checksum	boolean	GLOBAL
max_allowed_packet	integer	GLOBAL | SESSION
max_binlog_cache_size	integer	GLOBAL
max_binlog_size	integer	GLOBAL
max_binlog_stmt_cache_size	integer	GLOBAL
max_connect_errors	integer	GLOBAL
max_connections	integer	GLOBAL
max_delayed_threads	integer	GLOBAL | SESSION
max_error_count	integer	GLOBAL | SESSION
max_execution_time	integer	GLOBAL | SESSION
max_heap_table_size	integer	GLOBAL | SESSION
max_insert_delayed_threads	integer	GLOBAL | SESSION
max_join_size	integer	GLOBAL | SESSION
max_length_for_sort_data	integer	GLOBAL | SESSION
max_points_in_geometry	integer	GLOBAL
max_prepared_stmt_count	integer	GLOBAL
max_relay_log_size	integer	GLOBAL
max_seeks_for_key	integer	GLOBAL | SESSION
max_sort_length	integer	GLOBAL | SESSION
max_sp_recursion_depth	integer	GLOBAL | SESSION
max_tmp_tables	integer	GLOBAL | SESSION
max_user_connections	integer	GLOBAL | SESSION
max_write_lock_count	integer	GLOBAL
min_examined_row_limit	integer	GLOBAL | SESSION
myisam_data_pointer_size	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
myisam_max_sort_file_size	integer	GLOBAL
myisam_repair_threads	integer	GLOBAL | SESSION
myisam_sort_buffer_size	integer	GLOBAL | SESSION
myisam_stats_method	enumeration	GLOBAL | SESSION
myisam_use_mmap	boolean	GLOBAL
mysql_native_password_proxy_users	boolean	GLOBAL
mysqlx_connect_timeout	integer	GLOBAL
mysqlx_max_connections	integer	GLOBAL
ndb_blob_write_batch_bytes	integer	GLOBAL | SESSION
ndb_deferred_constraints	integer	GLOBAL | SESSION
ndb_deferred_constraints	integer	GLOBAL | SESSION
ndb_distribution	enumeration	GLOBAL
ndb_distribution={KEYHASH|LINHASH}	enumeration	GLOBAL
ndb_eventbuffer_free_percent	integer	GLOBAL
ndb_eventbuffer_max_alloc	integer	GLOBAL
ndb_force_send	boolean	GLOBAL | SESSION
ndb_index_stat_enable	boolean	GLOBAL | SESSION
ndb_index_stat_option	string	GLOBAL | SESSION
ndb_join_pushdown	boolean	GLOBAL | SESSION
ndb_log_binlog_index	boolean	GLOBAL
ndb_log_empty_epochs	boolean	GLOBAL
ndb_log_empty_update	boolean	GLOBAL
ndb_log_updated_only	boolean	GLOBAL
ndb_optimization_delay	integer	GLOBAL
ndb_recv_thread_cpu_mask	bitmap	GLOBAL
ndb_report_thresh_binlog_epoch_slip	integer	GLOBAL
ndb_report_thresh_binlog_mem_usage	integer	GLOBAL
ndb_show_foreign_key_mock_tables	boolean	GLOBAL
ndb_table_no_logging	boolean	SESSION
ndb_use_transactions	boolean	GLOBAL | SESSION
ndbinfo_max_rows	integer	GLOBAL | SESSION
ndbinfo_show_hidden	boolean	GLOBAL | SESSION
net_buffer_length	integer	GLOBAL | SESSION
net_read_timeout	integer	GLOBAL | SESSION
net_retry_count	integer	GLOBAL | SESSION
net_write_timeout	integer	GLOBAL | SESSION
new	boolean	GLOBAL | SESSION
offline_mode	boolean	GLOBAL
old_alter_table	boolean	GLOBAL | SESSION
old_passwords	enumeration	GLOBAL | SESSION
optimizer_prune_level	boolean	GLOBAL | SESSION

Variable Name	Variable Type	Variable Scope
optimizer_search_depth	integer	GLOBAL | SESSION
optimizer_switch	set	GLOBAL | SESSION
optimizer_trace	string	GLOBAL | SESSION
optimizer_trace_features	string	GLOBAL | SESSION
optimizer_trace_limit	integer	GLOBAL | SESSION
optimizer_trace_max_mem_size	integer	GLOBAL | SESSION
optimizer_trace_offset	integer	GLOBAL | SESSION
original_commit_timestamp	numeric	SESSION
parser_max_mem_size	integer	GLOBAL | SESSION
preload_buffer_size	integer	GLOBAL | SESSION
profiling	boolean	GLOBAL | SESSION
profiling_history_size	integer	GLOBAL | SESSION
pseudo_slave_mode	integer	SESSION
pseudo_thread_id	integer	SESSION
query_alloc_block_size	integer	GLOBAL | SESSION
query_cache_limit	integer	GLOBAL
query_cache_min_res_unit	integer	GLOBAL
query_cache_size	integer	GLOBAL
query_cache_type	enumeration	GLOBAL | SESSION
query_cache_wlock_invalidate	boolean	GLOBAL | SESSION
query_prealloc_size	integer	GLOBAL | SESSION
rand_seed1	integer	SESSION
rand_seed2	integer	SESSION
range_alloc_block_size	integer	GLOBAL | SESSION
range_optimizer_max_mem_size	integer	GLOBAL | SESSION
rbr_exec_mode	enumeration	SESSION
read_buffer_size	integer	GLOBAL | SESSION
read_only	boolean	GLOBAL
read_rnd_buffer_size	integer	GLOBAL | SESSION
relay_log_info_repository	string	GLOBAL
relay_log_purge	boolean	GLOBAL
require_secure_transport	boolean	GLOBAL
rewriter_enabled	boolean	GLOBAL
rewriter_verbose	integer	GLOBAL
rpl_semi_sync_master_enabled	boolean	GLOBAL
rpl_semi_sync_master_timeout	integer	GLOBAL
rpl_semi_sync_master_trace_level	integer	GLOBAL
rpl_semi_sync_master_wait_for_slave_count	integer	GLOBAL
rpl_semi_sync_master_wait_no_slave	boolean	GLOBAL
rpl_semi_sync_master_wait_point	enumeration	GLOBAL
rpl_semi_sync_slave_enabled	boolean	GLOBAL

Variable Name	Variable Type	Variable Scope
rpl_semi_sync_slave_trace_level	integer	GLOBAL
rpl_stop_slave_timeout	integer	GLOBAL
schema_definition_cache	integer	GLOBAL
secure_auth	boolean	GLOBAL
server_id	integer	GLOBAL
session_track_gtids	enumeration	GLOBAL | SESSION
session_track_schema	boolean	GLOBAL | SESSION
session_track_state_change	boolean	GLOBAL | SESSION
session_track_system_variables	string	GLOBAL | SESSION
sha256_password_proxy_users	boolean	GLOBAL
show_compatibility_56	boolean	GLOBAL
show_old_temporals	boolean	GLOBAL | SESSION
slave_allow_batching	boolean	GLOBAL
slave_checkpoint_group=#	integer	GLOBAL
slave_checkpoint_period=#	integer	GLOBAL
slave_compressed_protocol	boolean	GLOBAL
slave_exec_mode	enumeration	GLOBAL
slave_max_allowed_packet	integer	GLOBAL
slave_net_timeout	integer	GLOBAL
slave_parallel_type	enumeration	GLOBAL
slave_parallel_workers	integer	GLOBAL
slave_pending_jobs_size_max	integer	GLOBAL
slave_preserve_commit_order	boolean	GLOBAL
slave_rows_search_algorithms=list	set	GLOBAL
slave_sql_verify_checksum	boolean	GLOBAL
slave_transaction_retries	integer	GLOBAL
slow_launch_time	integer	GLOBAL
slow_query_log	boolean	GLOBAL
slow_query_log_file	filename	GLOBAL
sort_buffer_size	integer	GLOBAL | SESSION
sql_auto_is_null	boolean	GLOBAL | SESSION
sql_big_selects	boolean	GLOBAL | SESSION
sql_buffer_result	boolean	GLOBAL | SESSION
sql_log_bin	boolean	SESSION
sql_log_off	boolean	GLOBAL | SESSION
sql_mode	set	GLOBAL | SESSION
sql_notes	boolean	GLOBAL | SESSION
sql_quote_show_create	boolean	GLOBAL | SESSION
sql_safe_updates	boolean	GLOBAL | SESSION
sql_select_limit	integer	GLOBAL | SESSION
sql_slave_skip_counter	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
sql_warnings	boolean	GLOBAL | SESSION
stored_program_cache	integer	GLOBAL
stored_program_definition_cache	integer	GLOBAL
super_read_only	boolean	GLOBAL
sync_binlog	integer	GLOBAL
sync_master_info	integer	GLOBAL
sync_relay_log	integer	GLOBAL
sync_relay_log_info	integer	GLOBAL
table_definition_cache	integer	GLOBAL
table_open_cache	integer	GLOBAL
tablespace_definition_cache	integer	GLOBAL
temptable_max_ram	integer	GLOBAL
thread_cache_size	integer	GLOBAL
time_zone	string	GLOBAL | SESSION
timestamp	numeric	SESSION
tmp_table_size	integer	GLOBAL | SESSION
transaction_alloc_block_size	integer	GLOBAL | SESSION
transaction_allow_batching	boolean	SESSION
transaction_prealloc_size	integer	GLOBAL | SESSION
transaction_write_set_extraction	enumeration	GLOBAL | SESSION
tx_isolation	enumeration	GLOBAL | SESSION
tx_read_only	boolean	GLOBAL | SESSION
unique_checks	boolean	GLOBAL | SESSION
updatable_views_with_limit	boolean	GLOBAL | SESSION
validate_password_check_user_name	boolean	GLOBAL
validate_password_dictionary_file	filename	GLOBAL
validate_password_length	integer	GLOBAL
validate_password_mixed_case_count	integer	GLOBAL
validate_password_number_count	integer	GLOBAL
validate_password_policy	enumeration	GLOBAL
validate_password_special_char_count	integer	GLOBAL
version_tokens_session	string	GLOBAL | SESSION
wait_timeout	integer	GLOBAL | SESSION

5.1.7 Server Status Variables

The MySQL server maintains many status variables that provide information about its operation. You can view these variables and their values by using the `SHOW [GLOBAL | SESSION] STATUS` statement (see Section 13.7.5.35, “SHOW STATUS Syntax”). The optional `GLOBAL` keyword aggregates the values over all connections, and `SESSION` shows the values for the current connection.

```
mysql> SHOW GLOBAL STATUS;
+-----------------------------------+------------+
| Variable_name                     | Value      |
+-----------------------------------+------------+
```

```
| Aborted_clients                   | 0          |
| Aborted_connects                  | 0          |
| Bytes_received                    | 155372598  |
| Bytes_sent                        | 1176560426 |
...
| Connections                       | 30023      |
| Created_tmp_disk_tables           | 0          |
| Created_tmp_files                 | 3          |
| Created_tmp_tables                | 2          |
...
| Threads_created                   | 217        |
| Threads_running                   | 88         |
| Uptime                            | 1389872    |
+-----------------------------------+------------+
```

Several status variables provide statement counts. To determine the number of statements executed, use these relationships:

```
  SUM(Com_xxx) + Qcache_hits
= Questions + statements executed within stored programs
= Queries
```

Many status variables are reset to 0 by the `FLUSH STATUS` statement.

The following table lists all available server status variables:

Table 5.4 Status Variable Summary

Variable Name	Variable Type	Variable Scope
`Aborted_clients`	integer	`GLOBAL`
`Aborted_connects`	integer	`GLOBAL`
`Acl_cache_items_count`		`GLOBAL`
`Binlog_cache_disk_use`	integer	`GLOBAL`
`Binlog_cache_use`	integer	`GLOBAL`
`Binlog_stmt_cache_disk_use`	integer	`GLOBAL`
`Binlog_stmt_cache_use`	integer	`GLOBAL`
`Bytes_received`	integer	`GLOBAL` | `SESSION`
`Bytes_sent`	integer	`GLOBAL` | `SESSION`
`Com_admin_commands`	integer	`GLOBAL` | `SESSION`
`Com_alter_db`	integer	`GLOBAL` | `SESSION`
`Com_alter_event`	integer	`GLOBAL` | `SESSION`
`Com_alter_function`	integer	`GLOBAL` | `SESSION`
`Com_alter_procedure`	integer	`GLOBAL` | `SESSION`
`Com_alter_server`	integer	`GLOBAL` | `SESSION`
`Com_alter_table`	integer	`GLOBAL` | `SESSION`
`Com_alter_tablespace`	integer	`GLOBAL` | `SESSION`
`Com_alter_user`	integer	`GLOBAL` | `SESSION`
`Com_alter_user_default_role`		`GLOBAL`
`Com_analyze`	integer	`GLOBAL` | `SESSION`
`Com_assign_to_keycache`	integer	`GLOBAL` | `SESSION`
`Com_begin`	integer	`GLOBAL` | `SESSION`
`Com_binlog`	integer	`GLOBAL` | `SESSION`
`Com_call_procedure`	integer	`GLOBAL` | `SESSION`

Variable Name	Variable Type	Variable Scope
Com_change_db	integer	GLOBAL | SESSION
Com_change_master	integer	GLOBAL | SESSION
Com_change_repl_filter	integer	GLOBAL | SESSION
Com_check	integer	GLOBAL | SESSION
Com_checksum	integer	GLOBAL | SESSION
Com_commit	integer	GLOBAL | SESSION
Com_create_db	integer	GLOBAL | SESSION
Com_create_event	integer	GLOBAL | SESSION
Com_create_function	integer	GLOBAL | SESSION
Com_create_index	integer	GLOBAL | SESSION
Com_create_procedure	integer	GLOBAL | SESSION
Com_create_role		GLOBAL
Com_create_server	integer	GLOBAL | SESSION
Com_create_table	integer	GLOBAL | SESSION
Com_create_trigger	integer	GLOBAL | SESSION
Com_create_udf	integer	GLOBAL | SESSION
Com_create_user	integer	GLOBAL | SESSION
Com_create_view	integer	GLOBAL | SESSION
Com_dealloc_sql	integer	GLOBAL | SESSION
Com_delete	integer	GLOBAL | SESSION
Com_delete_multi	integer	GLOBAL | SESSION
Com_do	integer	GLOBAL | SESSION
Com_drop_db	integer	GLOBAL | SESSION
Com_drop_event	integer	GLOBAL | SESSION
Com_drop_function	integer	GLOBAL | SESSION
Com_drop_index	integer	GLOBAL | SESSION
Com_drop_procedure	integer	GLOBAL | SESSION
Com_drop_role		GLOBAL
Com_drop_server	integer	GLOBAL | SESSION
Com_drop_table	integer	GLOBAL | SESSION
Com_drop_trigger	integer	GLOBAL | SESSION
Com_drop_user	integer	GLOBAL | SESSION
Com_drop_view	integer	GLOBAL | SESSION
Com_empty_query	integer	GLOBAL | SESSION
Com_execute_sql	integer	GLOBAL | SESSION
Com_explain_other	integer	GLOBAL | SESSION
Com_flush	integer	GLOBAL | SESSION
Com_get_diagnostics	integer	GLOBAL | SESSION
Com_grant	integer	GLOBAL | SESSION
Com_grant_roles		GLOBAL
Com_group_replication_start		GLOBAL

Variable Name	Variable Type	Variable Scope
Com_group_replication_stop		GLOBAL
Com_ha_close	integer	GLOBAL | SESSION
Com_ha_open	integer	GLOBAL | SESSION
Com_ha_read	integer	GLOBAL | SESSION
Com_help	integer	GLOBAL | SESSION
Com_insert	integer	GLOBAL | SESSION
Com_insert_select	integer	GLOBAL | SESSION
Com_install_component		GLOBAL
Com_install_plugin	integer	GLOBAL | SESSION
Com_kill	integer	GLOBAL | SESSION
Com_load	integer	GLOBAL | SESSION
Com_lock_tables	integer	GLOBAL | SESSION
Com_optimize	integer	GLOBAL | SESSION
Com_preload_keys	integer	GLOBAL | SESSION
Com_prepare_sql	integer	GLOBAL | SESSION
Com_purge	integer	GLOBAL | SESSION
Com_purge_before_date	integer	GLOBAL | SESSION
Com_release_savepoint	integer	GLOBAL | SESSION
Com_rename_table	integer	GLOBAL | SESSION
Com_rename_user	integer	GLOBAL | SESSION
Com_repair	integer	GLOBAL | SESSION
Com_replace	integer	GLOBAL | SESSION
Com_replace_select	integer	GLOBAL | SESSION
Com_reset	integer	GLOBAL | SESSION
Com_resignal	integer	GLOBAL | SESSION
Com_revoke	integer	GLOBAL | SESSION
Com_revoke_all	integer	GLOBAL | SESSION
Com_revoke_roles		GLOBAL
Com_rollback	integer	GLOBAL | SESSION
Com_rollback_to_savepoint	integer	GLOBAL | SESSION
Com_savepoint	integer	GLOBAL | SESSION
Com_select	integer	GLOBAL | SESSION
Com_set_option	integer	GLOBAL | SESSION
Com_set_role		GLOBAL
Com_show_authors	integer	GLOBAL | SESSION
Com_show_binlog_events	integer	GLOBAL | SESSION
Com_show_binlogs	integer	GLOBAL | SESSION
Com_show_charsets	integer	GLOBAL | SESSION
Com_show_collations	integer	GLOBAL | SESSION
Com_show_contributors	integer	GLOBAL | SESSION
Com_show_create_db	integer	GLOBAL | SESSION

Variable Name	Variable Type	Variable Scope
Com_show_create_event	integer	GLOBAL | SESSION
Com_show_create_func	integer	GLOBAL | SESSION
Com_show_create_proc	integer	GLOBAL | SESSION
Com_show_create_table	integer	GLOBAL | SESSION
Com_show_create_trigger	integer	GLOBAL | SESSION
Com_show_create_user	integer	GLOBAL | SESSION
Com_show_databases	integer	GLOBAL | SESSION
Com_show_engine_logs	integer	GLOBAL | SESSION
Com_show_engine_mutex	integer	GLOBAL | SESSION
Com_show_engine_status	integer	GLOBAL | SESSION
Com_show_errors	integer	GLOBAL | SESSION
Com_show_events	integer	GLOBAL | SESSION
Com_show_fields	integer	GLOBAL | SESSION
Com_show_function_code	integer	GLOBAL | SESSION
Com_show_function_status	integer	GLOBAL | SESSION
Com_show_grants	integer	GLOBAL | SESSION
Com_show_keys	integer	GLOBAL | SESSION
Com_show_master_status	integer	GLOBAL | SESSION
Com_show_ndb_status	integer	GLOBAL | SESSION
Com_show_new_master	integer	GLOBAL | SESSION
Com_show_open_tables	integer	GLOBAL | SESSION
Com_show_plugins	integer	GLOBAL | SESSION
Com_show_privileges	integer	GLOBAL | SESSION
Com_show_procedure_code	integer	GLOBAL | SESSION
Com_show_procedure_status	integer	GLOBAL | SESSION
Com_show_processlist	integer	GLOBAL | SESSION
Com_show_profile	integer	GLOBAL | SESSION
Com_show_profiles	integer	GLOBAL | SESSION
Com_show_relaylog_events	integer	GLOBAL | SESSION
Com_show_slave_hosts	integer	GLOBAL | SESSION
Com_show_slave_status	integer	GLOBAL | SESSION
Com_show_slave_status_nonblocking	integer	GLOBAL | SESSION
Com_show_status	integer	GLOBAL | SESSION
Com_show_storage_engines	integer	GLOBAL | SESSION
Com_show_table_status	integer	GLOBAL | SESSION
Com_show_tables	integer	GLOBAL | SESSION
Com_show_triggers	integer	GLOBAL | SESSION
Com_show_variables	integer	GLOBAL | SESSION
Com_show_warnings	integer	GLOBAL | SESSION
Com_shutdown	integer	GLOBAL | SESSION
Com_signal	integer	GLOBAL | SESSION

Variable Name	Variable Type	Variable Scope
Com_slave_start	integer	GLOBAL | SESSION
Com_slave_stop	integer	GLOBAL | SESSION
Com_stmt_close	integer	GLOBAL | SESSION
Com_stmt_execute	integer	GLOBAL | SESSION
Com_stmt_fetch	integer	GLOBAL | SESSION
Com_stmt_prepare	integer	GLOBAL | SESSION
Com_stmt_reprepare	integer	GLOBAL | SESSION
Com_stmt_reset	integer	GLOBAL | SESSION
Com_stmt_send_long_data	integer	GLOBAL | SESSION
Com_truncate	integer	GLOBAL | SESSION
Com_uninstall_component		GLOBAL
Com_uninstall_plugin	integer	GLOBAL | SESSION
Com_unlock_tables	integer	GLOBAL | SESSION
Com_update	integer	GLOBAL | SESSION
Com_update_multi	integer	GLOBAL | SESSION
Com_xa_commit	integer	GLOBAL | SESSION
Com_xa_end	integer	GLOBAL | SESSION
Com_xa_prepare	integer	GLOBAL | SESSION
Com_xa_recover	integer	GLOBAL | SESSION
Com_xa_rollback	integer	GLOBAL | SESSION
Com_xa_start	integer	GLOBAL | SESSION
Compression	integer	SESSION
Connection_control_delay_generated		GLOBAL
Connection_errors_accept	integer	GLOBAL
Connection_errors_internal	integer	GLOBAL
Connection_errors_max_connections	integer	GLOBAL
Connection_errors_peer_addr	integer	GLOBAL
Connection_errors_select	integer	GLOBAL
Connection_errors_tcpwrap	integer	GLOBAL
Connections	integer	GLOBAL
Created_tmp_disk_tables	integer	GLOBAL | SESSION
Created_tmp_files	integer	GLOBAL
Created_tmp_tables	integer	GLOBAL | SESSION
Delayed_errors	integer	GLOBAL
Delayed_insert_threads	integer	GLOBAL
Delayed_writes	integer	GLOBAL
Flush_commands	integer	GLOBAL
Handler_commit	integer	GLOBAL | SESSION
Handler_delete	integer	GLOBAL | SESSION
Handler_external_lock	integer	GLOBAL | SESSION
Handler_mrr_init	integer	GLOBAL | SESSION

Variable Name	Variable Type	Variable Scope
Handler_prepare	integer	GLOBAL | SESSION
Handler_read_first	integer	GLOBAL | SESSION
Handler_read_key	integer	GLOBAL | SESSION
Handler_read_last	integer	GLOBAL | SESSION
Handler_read_next	integer	GLOBAL | SESSION
Handler_read_prev	integer	GLOBAL | SESSION
Handler_read_rnd	integer	GLOBAL | SESSION
Handler_read_rnd_next	integer	GLOBAL | SESSION
Handler_rollback	integer	GLOBAL | SESSION
Handler_savepoint	integer	GLOBAL | SESSION
Handler_savepoint_rollback	integer	GLOBAL | SESSION
Handler_update	integer	GLOBAL | SESSION
Handler_write	integer	GLOBAL | SESSION
Innodb_available_undo_logs	integer	GLOBAL
Innodb_buffer_pool_bytes_data	integer	GLOBAL
Innodb_buffer_pool_bytes_dirty	integer	GLOBAL
Innodb_buffer_pool_dump_status	string	GLOBAL
Innodb_buffer_pool_load_status	string	GLOBAL
Innodb_buffer_pool_pages_data	integer	GLOBAL
Innodb_buffer_pool_pages_dirty	integer	GLOBAL
Innodb_buffer_pool_pages_flushed	integer	GLOBAL
Innodb_buffer_pool_pages_free	integer	GLOBAL
Innodb_buffer_pool_pages_latched	integer	GLOBAL
Innodb_buffer_pool_pages_misc	integer	GLOBAL
Innodb_buffer_pool_pages_total	integer	GLOBAL
Innodb_buffer_pool_read_ahead	integer	GLOBAL
Innodb_buffer_pool_read_ahead_evicted	integer	GLOBAL
Innodb_buffer_pool_read_ahead_rnd	integer	GLOBAL
Innodb_buffer_pool_read_requests	integer	GLOBAL
Innodb_buffer_pool_reads	integer	GLOBAL
Innodb_buffer_pool_resize_status	string	GLOBAL
Innodb_buffer_pool_wait_free	integer	GLOBAL
Innodb_buffer_pool_write_requests	integer	GLOBAL
Innodb_data_fsyncs	integer	GLOBAL
Innodb_data_pending_fsyncs	integer	GLOBAL
Innodb_data_pending_reads	integer	GLOBAL
Innodb_data_pending_writes	integer	GLOBAL
Innodb_data_read	integer	GLOBAL
Innodb_data_reads	integer	GLOBAL
Innodb_data_writes	integer	GLOBAL
Innodb_data_written	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
Innodb_dblwr_pages_written	integer	GLOBAL
Innodb_dblwr_writes	integer	GLOBAL
Innodb_have_atomic_builtins	integer	GLOBAL
Innodb_log_waits	integer	GLOBAL
Innodb_log_write_requests	integer	GLOBAL
Innodb_log_writes	integer	GLOBAL
Innodb_num_open_files	integer	GLOBAL
Innodb_os_log_fsyncs	integer	GLOBAL
Innodb_os_log_pending_fsyncs	integer	GLOBAL
Innodb_os_log_pending_writes	integer	GLOBAL
Innodb_os_log_written	integer	GLOBAL
Innodb_page_size	integer	GLOBAL
Innodb_pages_created	integer	GLOBAL
Innodb_pages_read	integer	GLOBAL
Innodb_pages_written	integer	GLOBAL
Innodb_row_lock_current_waits	integer	GLOBAL
Innodb_row_lock_time	integer	GLOBAL
Innodb_row_lock_time_avg	integer	GLOBAL
Innodb_row_lock_time_max	integer	GLOBAL
Innodb_row_lock_waits	integer	GLOBAL
Innodb_rows_deleted	integer	GLOBAL
Innodb_rows_inserted	integer	GLOBAL
Innodb_rows_read	integer	GLOBAL
Innodb_rows_updated	integer	GLOBAL
Innodb_truncated_status_writes	integer	GLOBAL
Key_blocks_not_flushed	integer	GLOBAL
Key_blocks_unused	integer	GLOBAL
Key_blocks_used	integer	GLOBAL
Key_read_requests	integer	GLOBAL
Key_reads	integer	GLOBAL
Key_write_requests	integer	GLOBAL
Key_writes	integer	GLOBAL
Last_query_cost	numeric	SESSION
Last_query_partial_plans	integer	SESSION
Locked_connects	integer	GLOBAL
Max_execution_time_exceeded	integer	GLOBAL | SESSION
Max_execution_time_set	integer	GLOBAL | SESSION
Max_execution_time_set_failed	integer	GLOBAL | SESSION
Max_used_connections	integer	GLOBAL
Max_used_connections_time	datetime	GLOBAL
mecab_charset	string	GLOBAL

Variable Name	Variable Type	Variable Scope
Mysqlx_address	string	GLOBAL
Mysqlx_bytes_received	integer	GLOBAL | SESSION
Mysqlx_bytes_sent	integer	GLOBAL | SESSION
Mysqlx_connection_accept_errors	integer	GLOBAL | SESSION
Mysqlx_connection_errors	integer	GLOBAL | SESSION
Mysqlx_connections_accepted	integer	GLOBAL
Mysqlx_connections_closed	integer	GLOBAL
Mysqlx_connections_rejected	integer	GLOBAL
Mysqlx_crud_create_view	integer	GLOBAL | SESSION
Mysqlx_crud_delete	integer	GLOBAL | SESSION
Mysqlx_crud_drop_view	integer	GLOBAL | SESSION
Mysqlx_crud_find	integer	GLOBAL | SESSION
Mysqlx_crud_insert	integer	GLOBAL | SESSION
Mysqlx_crud_modify_view	integer	GLOBAL | SESSION
Mysqlx_crud_update	integer	GLOBAL | SESSION
Mysqlx_errors_sent	integer	GLOBAL | SESSION
Mysqlx_errors_unknown_message_type	integer	GLOBAL | SESSION
Mysqlx_expect_close	integer	GLOBAL | SESSION
Mysqlx_expect_open	integer	GLOBAL | SESSION
Mysqlx_init_error	integer	GLOBAL | SESSION
Mysqlx_notice_other_sent	integer	GLOBAL | SESSION
Mysqlx_notice_warning_sent	integer	GLOBAL | SESSION
Mysqlx_port	string	GLOBAL
Mysqlx_rows_sent	integer	GLOBAL | SESSION
Mysqlx_sessions	integer	GLOBAL
Mysqlx_sessions_accepted	integer	GLOBAL
Mysqlx_sessions_closed	integer	GLOBAL
Mysqlx_sessions_fatal_error	integer	GLOBAL
Mysqlx_sessions_killed	integer	GLOBAL
Mysqlx_sessions_rejected	integer	GLOBAL
Mysqlx_socket	string	GLOBAL
Mysqlx_ssl_accept_renegotiates	integer	GLOBAL
Mysqlx_ssl_accepts	integer	GLOBAL
Mysqlx_ssl_active	integer	GLOBAL | SESSION
Mysqlx_ssl_cipher	integer	GLOBAL | SESSION
Mysqlx_ssl_cipher_list	integer	GLOBAL | SESSION
Mysqlx_ssl_ctx_verify_depth	integer	GLOBAL | SESSION
Mysqlx_ssl_ctx_verify_mode	integer	GLOBAL | SESSION
Mysqlx_ssl_finished_accepts	integer	GLOBAL
Mysqlx_ssl_server_not_after	integer	GLOBAL
Mysqlx_ssl_server_not_before	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
Mysqlx_ssl_verify_depth	integer	GLOBAL
Mysqlx_ssl_verify_mode	integer	GLOBAL
Mysqlx_ssl_version	integer	GLOBAL | SESSION
Mysqlx_stmt_create_collection	integer	GLOBAL | SESSION
Mysqlx_stmt_create_collection_index	integer	GLOBAL | SESSION
Mysqlx_stmt_disable_notices	integer	GLOBAL | SESSION
Mysqlx_stmt_drop_collection	integer	GLOBAL | SESSION
Mysqlx_stmt_drop_collection_index	integer	GLOBAL | SESSION
Mysqlx_stmt_enable_notices	integer	GLOBAL | SESSION
Mysqlx_stmt_ensure_collection	string	GLOBAL | SESSION
Mysqlx_stmt_execute_mysqlx	integer	GLOBAL | SESSION
Mysqlx_stmt_execute_sql	integer	GLOBAL | SESSION
Mysqlx_stmt_execute_xplugin	integer	GLOBAL | SESSION
Mysqlx_stmt_kill_client	integer	GLOBAL | SESSION
Mysqlx_stmt_list_clients	integer	GLOBAL | SESSION
Mysqlx_stmt_list_notices	integer	GLOBAL | SESSION
Mysqlx_stmt_list_objects	integer	GLOBAL | SESSION
Mysqlx_stmt_ping	integer	GLOBAL | SESSION
Mysqlx_worker_threads	integer	GLOBAL
Mysqlx_worker_threads_active	integer	GLOBAL
Ndb_api_bytes_received_count	integer	GLOBAL
Ndb_api_bytes_received_count_session	integer	SESSION
Ndb_api_bytes_received_count_slave	integer	GLOBAL
Ndb_api_bytes_sent_count	integer	GLOBAL
Ndb_api_bytes_sent_count_slave	integer	GLOBAL
Ndb_api_event_bytes_count_injector	integer	GLOBAL
Ndb_api_event_data_count_injector	integer	GLOBAL
Ndb_api_event_nondata_count_injector	integer	GLOBAL
Ndb_api_pk_op_count	integer	GLOBAL
Ndb_api_pk_op_count_session	integer	SESSION
Ndb_api_pk_op_count_slave	integer	GLOBAL
Ndb_api_pruned_scan_count	integer	GLOBAL
Ndb_api_pruned_scan_count_session	integer	SESSION
Ndb_api_range_scan_count_slave	integer	GLOBAL
Ndb_api_read_row_count	integer	GLOBAL
Ndb_api_read_row_count_session	integer	SESSION
Ndb_api_scan_batch_count_slave	integer	GLOBAL
Ndb_api_table_scan_count	integer	GLOBAL
Ndb_api_table_scan_count_session	integer	SESSION
Ndb_api_trans_abort_count	integer	GLOBAL
Ndb_api_trans_abort_count_session	integer	SESSION

Variable Name	Variable Type	Variable Scope
Ndb_api_trans_abort_count_slave	integer	GLOBAL
Ndb_api_trans_close_count	integer	GLOBAL
Ndb_api_trans_close_count_session	integer	SESSION
Ndb_api_trans_close_count_slave	integer	GLOBAL
Ndb_api_trans_commit_count	integer	GLOBAL
Ndb_api_trans_commit_count_session	integer	SESSION
Ndb_api_trans_commit_count_slave	integer	GLOBAL
Ndb_api_trans_local_read_row_count_slave	integer	GLOBAL
Ndb_api_trans_start_count	integer	GLOBAL
Ndb_api_trans_start_count_session	integer	SESSION
Ndb_api_trans_start_count_slave	integer	GLOBAL
Ndb_api_uk_op_count	integer	GLOBAL
Ndb_api_uk_op_count_slave	integer	GLOBAL
Ndb_api_wait_exec_complete_count	integer	GLOBAL
Ndb_api_wait_exec_complete_count_session	integer	SESSION
Ndb_api_wait_exec_complete_count_slave	integer	GLOBAL
Ndb_api_wait_meta_request_count	integer	GLOBAL
Ndb_api_wait_meta_request_count_session	integer	SESSION
Ndb_api_wait_nanos_count	integer	GLOBAL
Ndb_api_wait_nanos_count_session	integer	SESSION
Ndb_api_wait_nanos_count_slave	integer	GLOBAL
Ndb_api_wait_scan_result_count	integer	GLOBAL
Ndb_api_wait_scan_result_count_session	integer	SESSION
Ndb_api_wait_scan_result_count_slave	integer	GLOBAL
Ndb_cluster_node_id	integer	GLOBAL | SESSION
Ndb_config_from_host	integer	GLOBAL | SESSION
Ndb_config_from_port	integer	GLOBAL | SESSION
Ndb_conflict_fn_epoch_trans	integer	GLOBAL
Ndb_conflict_fn_max	integer	GLOBAL
Ndb_conflict_fn_old	integer	GLOBAL
Ndb_conflict_trans_detect_iter_count	integer	GLOBAL
Ndb_conflict_trans_row_reject_count	integer	GLOBAL
Ndb_last_commit_epoch_server	integer	GLOBAL
Ndb_last_commit_epoch_session	integer	SESSION
Ndb_number_of_data_nodes	integer	GLOBAL
Ndb_pushed_queries_defined	integer	GLOBAL
Ndb_pushed_queries_executed	integer	GLOBAL
Ndb_scan_count	integer	GLOBAL
Not_flushed_delayed_rows	integer	GLOBAL
Ongoing_anonymous_gtid_violating_transaction_count	integer	GLOBAL
Ongoing_anonymous_transaction_count	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
Ongoing_automatic_gtid_violating_transaction_count	integer	GLOBAL
Open_files	integer	GLOBAL
Open_streams	integer	GLOBAL
Open_table_definitions	integer	GLOBAL
Open_tables	integer	GLOBAL | SESSION
Opened_files	integer	GLOBAL
Opened_table_definitions	integer	GLOBAL | SESSION
Opened_tables	integer	GLOBAL | SESSION
original_commit_timestamp	numeric	SESSION
Performance_schema_accounts_lost	integer	GLOBAL
Performance_schema_cond_classes_lost	integer	GLOBAL
Performance_schema_cond_instances_lost	integer	GLOBAL
Performance_schema_digest_lost	integer	GLOBAL
Performance_schema_file_classes_lost	integer	GLOBAL
Performance_schema_file_handles_lost	integer	GLOBAL
Performance_schema_file_instances_lost	integer	GLOBAL
Performance_schema_hosts_lost	integer	GLOBAL
Performance_schema_index_stat_lost	integer	GLOBAL
Performance_schema_locker_lost	integer	GLOBAL
Performance_schema_memory_classes_lost	integer	GLOBAL
Performance_schema_metadata_lock_lost	integer	GLOBAL
Performance_schema_mutex_classes_lost	integer	GLOBAL
Performance_schema_mutex_instances_lost	integer	GLOBAL
Performance_schema_nested_statement_lost	integer	GLOBAL
Performance_schema_prepared_statements_lost	integer	GLOBAL
Performance_schema_program_lost	integer	GLOBAL
Performance_schema_rwlock_classes_lost	integer	GLOBAL
Performance_schema_rwlock_instances_lost	integer	GLOBAL
Performance_schema_session_connect_attrs_longest_seen	integer	GLOBAL
Performance_schema_session_connect_attrs_lost	integer	GLOBAL
Performance_schema_socket_classes_lost	integer	GLOBAL
Performance_schema_socket_instances_lost	integer	GLOBAL
Performance_schema_stage_classes_lost	integer	GLOBAL
Performance_schema_statement_classes_lost	integer	GLOBAL
Performance_schema_table_handles_lost	integer	GLOBAL
Performance_schema_table_instances_lost	integer	GLOBAL
Performance_schema_table_lock_stat_lost	integer	GLOBAL
Performance_schema_thread_classes_lost	integer	GLOBAL
Performance_schema_thread_instances_lost	integer	GLOBAL
Performance_schema_users_lost	integer	GLOBAL
Prepared_stmt_count	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
Qcache_free_blocks	integer	GLOBAL
Qcache_free_memory	integer	GLOBAL
Qcache_hits	integer	GLOBAL
Qcache_inserts	integer	GLOBAL
Qcache_lowmem_prunes	integer	GLOBAL
Qcache_not_cached	integer	GLOBAL
Qcache_queries_in_cache	integer	GLOBAL
Qcache_total_blocks	integer	GLOBAL
Queries	integer	GLOBAL | SESSION
Questions	integer	GLOBAL | SESSION
Rewriter_number_loaded_rules	integer	GLOBAL
Rewriter_number_reloads	integer	GLOBAL
Rewriter_number_rewritten_queries	integer	GLOBAL
Rewriter_reload_error	boolean	GLOBAL
Rpl_semi_sync_master_clients	integer	GLOBAL
Rpl_semi_sync_master_net_avg_wait_time	integer	GLOBAL
Rpl_semi_sync_master_net_wait_time	integer	GLOBAL
Rpl_semi_sync_master_net_waits	integer	GLOBAL
Rpl_semi_sync_master_no_times	integer	GLOBAL
Rpl_semi_sync_master_no_tx	integer	GLOBAL
Rpl_semi_sync_master_status	boolean	GLOBAL
Rpl_semi_sync_master_timefunc_failures	integer	GLOBAL
Rpl_semi_sync_master_tx_avg_wait_time	integer	GLOBAL
Rpl_semi_sync_master_tx_wait_time	integer	GLOBAL
Rpl_semi_sync_master_tx_waits	integer	GLOBAL
Rpl_semi_sync_master_wait_pos_backtraverse	integer	GLOBAL
Rpl_semi_sync_master_wait_sessions	integer	GLOBAL
Rpl_semi_sync_master_yes_tx	integer	GLOBAL
Rpl_semi_sync_slave_status	boolean	GLOBAL
Rsa_public_key	string	GLOBAL
Select_full_join	integer	GLOBAL | SESSION
Select_full_range_join	integer	GLOBAL | SESSION
Select_range	integer	GLOBAL | SESSION
Select_range_check	integer	GLOBAL | SESSION
Select_scan	integer	GLOBAL | SESSION
Slave_heartbeat_period	numeric	GLOBAL
Slave_last_heartbeat		GLOBAL
Slave_open_temp_tables	integer	GLOBAL
Slave_received_heartbeats		GLOBAL
Slave_retried_transactions	integer	GLOBAL
Slave_running		GLOBAL

Variable Name	Variable Type	Variable Scope
Slow_launch_threads	integer	GLOBAL | SESSION
Slow_queries	integer	GLOBAL | SESSION
Sort_merge_passes	integer	GLOBAL | SESSION
Sort_range	integer	GLOBAL | SESSION
Sort_rows	integer	GLOBAL | SESSION
Sort_scan	integer	GLOBAL | SESSION
Ssl_accept_renegotiates	integer	GLOBAL
Ssl_accepts	integer	GLOBAL
Ssl_callback_cache_hits	integer	GLOBAL
Ssl_cipher	string	GLOBAL | SESSION
Ssl_cipher_list	string	GLOBAL | SESSION
Ssl_client_connects	integer	GLOBAL
Ssl_connect_renegotiates	integer	GLOBAL
Ssl_ctx_verify_depth	integer	GLOBAL
Ssl_ctx_verify_mode	integer	GLOBAL
Ssl_default_timeout	integer	GLOBAL | SESSION
Ssl_finished_accepts	integer	GLOBAL
Ssl_finished_connects	integer	GLOBAL
Ssl_server_not_after	integer	GLOBAL | SESSION
Ssl_server_not_before	integer	GLOBAL | SESSION
Ssl_session_cache_hits	integer	GLOBAL
Ssl_session_cache_misses	integer	GLOBAL
Ssl_session_cache_mode	string	GLOBAL
Ssl_session_cache_overflows	integer	GLOBAL
Ssl_session_cache_size	integer	GLOBAL
Ssl_session_cache_timeouts	integer	GLOBAL
Ssl_sessions_reused	integer	GLOBAL | SESSION
Ssl_used_session_cache_entries	integer	GLOBAL
Ssl_verify_depth	integer	GLOBAL | SESSION
Ssl_verify_mode	integer	GLOBAL | SESSION
Ssl_version	string	GLOBAL | SESSION
Table_locks_immediate	integer	GLOBAL
Table_locks_waited	integer	GLOBAL
Table_open_cache_hits	integer	GLOBAL | SESSION
Table_open_cache_misses	integer	GLOBAL | SESSION
Table_open_cache_overflows	integer	GLOBAL | SESSION
Tc_log_max_pages_used	integer	GLOBAL
Tc_log_page_size	integer	GLOBAL
Tc_log_page_waits	integer	GLOBAL
Threads_cached	integer	GLOBAL
Threads_connected	integer	GLOBAL

Variable Name	Variable Type	Variable Scope
Threads_created	integer	GLOBAL
Threads_running	integer	GLOBAL
Uptime	integer	GLOBAL
Uptime_since_flush_status	integer	GLOBAL
validate_password_dictionary_file_last_parsed	datetime	GLOBAL
validate_password_dictionary_file_words_count	integer	GLOBAL

The status variables have the following meanings.

- Aborted_clients

  The number of connections that were aborted because the client died without closing the connection properly. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

- Aborted_connects

  The number of failed attempts to connect to the MySQL server. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

  For additional connection-related information, check the Connection_errors_*xxx* status variables and the host_cache table.

- Binlog_cache_disk_use

  The number of transactions that used the temporary binary log cache but that exceeded the value of binlog_cache_size and used a temporary file to store statements from the transaction.

  The number of nontransactional statements that caused the binary log transaction cache to be written to disk is tracked separately in the Binlog_stmt_cache_disk_use status variable.

- Acl_cache_items_count

  The number of cached privilege objects. Each object is the privilege combination of a user and its active roles.

- Binlog_cache_use

  The number of transactions that used the binary log cache.

- Binlog_stmt_cache_disk_use

  The number of nontransaction statements that used the binary log statement cache but that exceeded the value of binlog_stmt_cache_size and used a temporary file to store those statements.

- Binlog_stmt_cache_use

  The number of nontransactional statements that used the binary log statement cache.

- Bytes_received

  The number of bytes received from all clients.

- Bytes_sent

  The number of bytes sent to all clients.

- Com_*xxx*

The `Com_xxx` statement counter variables indicate the number of times each *`xxx`* statement has been executed. There is one status variable for each type of statement. For example, `Com_delete` and `Com_update` count `DELETE` and `UPDATE` statements, respectively. `Com_delete_multi` and `Com_update_multi` are similar but apply to `DELETE` and `UPDATE` statements that use multiple-table syntax.

If a query result is returned from query cache, the server increments the `Qcache_hits` status variable, not `Com_select`. See Section 8.10.3.4, “Query Cache Status and Maintenance”.

The discussion at the beginning of this section indicates how to relate these statement-counting status variables to other such variables.

All of the `Com_stmt_xxx` variables are increased even if a prepared statement argument is unknown or an error occurred during execution. In other words, their values correspond to the number of requests issued, not to the number of requests successfully completed.

The `Com_stmt_xxx` status variables are as follows:

- `Com_stmt_prepare`
- `Com_stmt_execute`
- `Com_stmt_fetch`
- `Com_stmt_send_long_data`
- `Com_stmt_reset`
- `Com_stmt_close`

Those variables stand for prepared statement commands. Their names refer to the `COM_xxx` command set used in the network layer. In other words, their values increase whenever prepared statement API calls such as `mysql_stmt_prepare()`, `mysql_stmt_execute()`, and so forth are executed. However, `Com_stmt_prepare`, `Com_stmt_execute` and `Com_stmt_close` also increase for `PREPARE`, `EXECUTE`, or `DEALLOCATE PREPARE`, respectively. Additionally, the values of the older statement counter variables `Com_prepare_sql`, `Com_execute_sql`, and `Com_dealloc_sql` increase for the `PREPARE`, `EXECUTE`, and `DEALLOCATE PREPARE` statements. `Com_stmt_fetch` stands for the total number of network round-trips issued when fetching from cursors.

`Com_stmt_reprepare` indicates the number of times statements were automatically reprepared by the server after metadata changes to tables or views referred to by the statement. A reprepare operation increments `Com_stmt_reprepare`, and also `Com_stmt_prepare`.

`Com_explain_other` indicates the number of `EXPLAIN FOR CONNECTION` statements executed. See Section 8.8.4, “Obtaining Execution Plan Information for a Named Connection”.

`Com_change_repl_filter` indicates the number of `CHANGE REPLICATION FILTER` statements executed.

- `Compression`

  Whether the client connection uses compression in the client/server protocol.

- `Connection_errors_xxx`

  These variables provide information about errors that occur during the client connection process. They are global only and represent error counts aggregated across connections from all hosts. These variables track errors not accounted for by the host cache (see Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”), such as errors that are not associated with TCP

connections, occur very early in the connection process (even before an IP address is known), or are not specific to any particular IP address (such as out-of-memory conditions).

- `Connection_errors_accept`

  The number of errors that occurred during calls to `accept()` on the listening port.

- `Connection_errors_internal`

  The number of connections refused due to internal errors in the server, such as failure to start a new thread or an out-of-memory condition.

- `Connection_errors_max_connections`

  The number of connections refused because the server `max_connections` limit was reached.

- `Connection_errors_peer_addr`

  The number of errors that occurred while searching for connecting client IP addresses.

- `Connection_errors_select`

  The number of errors that occurred during calls to `select()` or `poll()` on the listening port. (Failure of this operation does not necessarily means a client connection was rejected.)

- `Connection_errors_tcpwrap`

  The number of connections refused by the `libwrap` library.

- `Connections`

  The number of connection attempts (successful or not) to the MySQL server.

- `Created_tmp_disk_tables`

  The number of internal on-disk temporary tables created by the server while executing statements.

  If an internal temporary table is created initially as an in-memory table but becomes too large, MySQL automatically converts it to an on-disk table. The maximum size for in-memory temporary tables is the minimum of the `tmp_table_size` and `max_heap_table_size` values. If `Created_tmp_disk_tables` is large, you may want to increase the `tmp_table_size` or `max_heap_table_size` value to lessen the likelihood that internal temporary tables in memory will be converted to on-disk tables.

  You can compare the number of internal on-disk temporary tables created to the total number of internal temporary tables created by comparing the values of the `Created_tmp_disk_tables` and `Created_tmp_tables` variables.

  See also Section 8.4.4, “Internal Temporary Table Use in MySQL”.

- `Created_tmp_files`

  How many temporary files `mysqld` has created.

- `Created_tmp_tables`

  The number of internal temporary tables created by the server while executing statements.

  You can compare the number of internal on-disk temporary tables created to the total number of internal temporary tables created by comparing the values of the `Created_tmp_disk_tables` and `Created_tmp_tables` variables.

See also Section 8.4.4, “Internal Temporary Table Use in MySQL”.

Each invocation of the `SHOW STATUS` statement uses an internal temporary table and increments the global `Created_tmp_tables` value.

- `Delayed_errors`

  This status variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `Delayed_insert_threads`

  This status variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `Delayed_writes`

  This status variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `Flush_commands`

  The number of times the server flushes tables, whether because a user executed a `FLUSH TABLES` statement or due to internal server operation. It is also incremented by receipt of a `COM_REFRESH` packet. This is in contrast to `Com_flush`, which indicates how many `FLUSH` statements have been executed, whether `FLUSH TABLES`, `FLUSH LOGS`, and so forth.

- `Handler_commit`

  The number of internal `COMMIT` statements.

- `Handler_delete`

  The number of times that rows have been deleted from tables.

- `Handler_external_lock`

  The server increments this variable for each call to its `external_lock()` function, which generally occurs at the beginning and end of access to a table instance. There might be differences among storage engines. This variable can be used, for example, to discover for a statement that accesses a partitioned table how many partitions were pruned before locking occurred: Check how much the counter increased for the statement, subtract 2 (2 calls for the table itself), then divide by 2 to get the number of partitions locked.

- `Handler_mrr_init`

  The number of times the server uses a storage engine's own Multi-Range Read implementation for table access.

- `Handler_prepare`

  A counter for the prepare phase of two-phase commit operations.

- `Handler_read_first`

  The number of times the first entry in an index was read. If this value is high, it suggests that the server is doing a lot of full index scans; for example, `SELECT col1 FROM foo`, assuming that `col1` is indexed.

- `Handler_read_key`

The number of requests to read a row based on a key. If this value is high, it is a good indication that your tables are properly indexed for your queries.

- `Handler_read_last`

  The number of requests to read the last key in an index. With `ORDER BY`, the server will issue a first-key request followed by several next-key requests, whereas with `ORDER BY DESC`, the server will issue a last-key request followed by several previous-key requests.

- `Handler_read_next`

  The number of requests to read the next row in key order. This value is incremented if you are querying an index column with a range constraint or if you are doing an index scan.

- `Handler_read_prev`

  The number of requests to read the previous row in key order. This read method is mainly used to optimize `ORDER BY ... DESC`.

- `Handler_read_rnd`

  The number of requests to read a row based on a fixed position. This value is high if you are doing a lot of queries that require sorting of the result. You probably have a lot of queries that require MySQL to scan entire tables or you have joins that do not use keys properly.

- `Handler_read_rnd_next`

  The number of requests to read the next row in the data file. This value is high if you are doing a lot of table scans. Generally this suggests that your tables are not properly indexed or that your queries are not written to take advantage of the indexes you have.

- `Handler_rollback`

  The number of requests for a storage engine to perform a rollback operation.

- `Handler_savepoint`

  The number of requests for a storage engine to place a savepoint.

- `Handler_savepoint_rollback`

  The number of requests for a storage engine to roll back to a savepoint.

- `Handler_update`

  The number of requests to update a row in a table.

- `Handler_write`

  The number of requests to insert a row in a table.

- `Innodb_available_undo_logs`

  Note

  `Innodb_available_undo_logs` is removed in MySQL 8.0.2. As of MySQL 8.0.2, the number of available rollback segments per tablespace may be retrieved using `SHOW VARIABLES LIKE 'innodb_rollback_segments';`

  The total number of available `InnoDB` rollback segments. Supplements the `innodb_rollback_segments` system variable, which defines the number of active rollback segments.

One rollback segment always resides in the system tablespace, and 32 rollback segments are reserved for use by temporary tables and are hosted in the temporary tablespace (`ibtmp1`). See Section 15.4.11.1, “Temporary Table Undo Logs”.

If you initiate a MySQL instance with 32 or fewer rollback segments, `InnoDB` still assigns one rollback segment to the system tablespace and 32 rollback segments to the temporary tablespace. In this case, `Innodb_available_undo_logs` reports 33 available rollback segments even though the instance was initialized with a lesser `innodb_rollback_segments` value.

- `Innodb_buffer_pool_dump_status`

  The progress of an operation to record the pages held in the `InnoDB` buffer pool, triggered by the setting of `innodb_buffer_pool_dump_at_shutdown` or `innodb_buffer_pool_dump_now`.

  For related information and examples, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `Innodb_buffer_pool_load_status`

  The progress of an operation to warm up the `InnoDB` buffer pool by reading in a set of pages corresponding to an earlier point in time, triggered by the setting of `innodb_buffer_pool_load_at_startup` or `innodb_buffer_pool_load_now`. If the operation introduces too much overhead, you can cancel it by setting `innodb_buffer_pool_load_abort`.

  For related information and examples, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `Innodb_buffer_pool_bytes_data`

  The total number of bytes in the `InnoDB` buffer pool containing data. The number includes both dirty and clean pages. For more accurate memory usage calculations than with `Innodb_buffer_pool_pages_data`, when compressed tables cause the buffer pool to hold pages of different sizes.

- `Innodb_buffer_pool_pages_data`

  The number of pages in the `InnoDB` buffer pool containing data. The number includes both dirty and clean pages. When using compressed tables, the reported `Innodb_buffer_pool_pages_data` value may be larger than `Innodb_buffer_pool_pages_total` (Bug #59550).

- `Innodb_buffer_pool_bytes_dirty`

  The total current number of bytes held in dirty pages in the `InnoDB` buffer pool. For more accurate memory usage calculations than with `Innodb_buffer_pool_pages_dirty`, when compressed tables cause the buffer pool to hold pages of different sizes.

- `Innodb_buffer_pool_pages_dirty`

  The current number of dirty pages in the `InnoDB` buffer pool.

- `Innodb_buffer_pool_pages_flushed`

  The number of requests to flush pages from the `InnoDB` buffer pool.

- `Innodb_buffer_pool_pages_free`

  The number of free pages in the `InnoDB` buffer pool.

- `Innodb_buffer_pool_pages_latched`

The number of latched pages in the InnoDB buffer pool. These are pages currently being read or written, or that cannot be flushed or removed for some other reason. Calculation of this variable is expensive, so it is available only when the UNIV_DEBUG system is defined at server build time.

- Innodb_buffer_pool_pages_misc

  The number of pages in the InnoDB buffer pool that are busy because they have been allocated for administrative overhead, such as row locks or the adaptive hash index. This value can also be calculated as Innodb_buffer_pool_pages_total − Innodb_buffer_pool_pages_free − Innodb_buffer_pool_pages_data. When using compressed tables, Innodb_buffer_pool_pages_misc may report an out-of-bounds value (Bug #59550).

- Innodb_buffer_pool_pages_total

  The total size of the InnoDB buffer pool, in pages. When using compressed tables, the reported Innodb_buffer_pool_pages_data value may be larger than Innodb_buffer_pool_pages_total (Bug #59550)

- Innodb_buffer_pool_read_ahead

  The number of pages read into the InnoDB buffer pool by the read-ahead background thread.

- Innodb_buffer_pool_read_ahead_evicted

  The number of pages read into the InnoDB buffer pool by the read-ahead background thread that were subsequently evicted without having been accessed by queries.

- Innodb_buffer_pool_read_ahead_rnd

  The number of “random” read-aheads initiated by InnoDB. This happens when a query scans a large portion of a table but in random order.

- Innodb_buffer_pool_read_requests

  The number of logical read requests.

- Innodb_buffer_pool_reads

  The number of logical reads that InnoDB could not satisfy from the buffer pool, and had to read directly from disk.

- Innodb_buffer_pool_resize_status

  The status of an operation to resize the InnoDB buffer pool dynamically, triggered by setting the innodb_buffer_pool_size parameter dynamically. The innodb_buffer_pool_size parameter is dynamic, which allows you to resize the buffer pool without restarting the server. See Configuring InnoDB Buffer Pool Size Online for related information.

- Innodb_buffer_pool_wait_free

  Normally, writes to the InnoDB buffer pool happen in the background. When InnoDB needs to read or create a page and no clean pages are available, InnoDB flushes some dirty pages first and waits for that operation to finish. This counter counts instances of these waits. If innodb_buffer_pool_size has been set properly, this value should be small.

- Innodb_buffer_pool_write_requests

  The number of writes done to the InnoDB buffer pool.

- Innodb_data_fsyncs

The number of `fsync()` operations so far. The frequency of `fsync()` calls is influenced by the setting of the `innodb_flush_method` configuration option.

- `Innodb_data_pending_fsyncs`

  The current number of pending `fsync()` operations. The frequency of `fsync()` calls is influenced by the setting of the `innodb_flush_method` configuration option.

- `Innodb_data_pending_reads`

  The current number of pending reads.

- `Innodb_data_pending_writes`

  The current number of pending writes.

- `Innodb_data_read`

  The amount of data read since the server was started (in bytes).

- `Innodb_data_reads`

  The total number of data reads (OS file reads).

- `Innodb_data_writes`

  The total number of data writes.

- `Innodb_data_written`

  The amount of data written so far, in bytes.

- `Innodb_dblwr_pages_written`

  The number of pages that have been written to the doublewrite buffer. See Section 15.11.1, “InnoDB Disk I/O”.

- `Innodb_dblwr_writes`

  The number of doublewrite operations that have been performed. See Section 15.11.1, “InnoDB Disk I/O”.

- `Innodb_have_atomic_builtins`

  Indicates whether the server was built with atomic instructions.

- `Innodb_log_waits`

  The number of times that the log buffer was too small and a wait was required for it to be flushed before continuing.

- `Innodb_log_write_requests`

  The number of write requests for the `InnoDB` redo log.

- `Innodb_log_writes`

  The number of physical writes to the `InnoDB` redo log file.

- `Innodb_num_open_files`

  The number of files `InnoDB` currently holds open.

- `Innodb_os_log_fsyncs`

The number of `fsync()` writes done to the `InnoDB` redo log files.

- `Innodb_os_log_pending_fsyncs`

  The number of pending `fsync()` operations for the `InnoDB` redo log files.

- `Innodb_os_log_pending_writes`

  The number of pending writes to the `InnoDB` redo log files.

- `Innodb_os_log_written`

  The number of bytes written to the `InnoDB` redo log files.

- `Innodb_page_size`

  `InnoDB` page size (default 16KB). Many values are counted in pages; the page size enables them to be easily converted to bytes.

- `Innodb_pages_created`

  The number of pages created by operations on `InnoDB` tables.

- `Innodb_pages_read`

  The number of pages read from the `InnoDB` buffer pool by operations on `InnoDB` tables.

- `Innodb_pages_written`

  The number of pages written by operations on `InnoDB` tables.

- `Innodb_row_lock_current_waits`

  The number of row locks currently being waited for by operations on `InnoDB` tables.

- `Innodb_row_lock_time`

  The total time spent in acquiring row locks for `InnoDB` tables, in milliseconds.

- `Innodb_row_lock_time_avg`

  The average time to acquire a row lock for `InnoDB` tables, in milliseconds.

- `Innodb_row_lock_time_max`

  The maximum time to acquire a row lock for `InnoDB` tables, in milliseconds.

- `Innodb_row_lock_waits`

  The number of times operations on `InnoDB` tables had to wait for a row lock.

- `Innodb_rows_deleted`

  The number of rows deleted from `InnoDB` tables.

- `Innodb_rows_inserted`

  The number of rows inserted into `InnoDB` tables.

- `Innodb_rows_read`

  The number of rows read from `InnoDB` tables.

- `Innodb_rows_updated`

The number of rows updated in `InnoDB` tables.

- `Innodb_truncated_status_writes`

  The number of times output from the `SHOW ENGINE INNODB STATUS` statement has been truncated.

- `Key_blocks_not_flushed`

  The number of key blocks in the `MyISAM` key cache that have changed but have not yet been flushed to disk.

- `Key_blocks_unused`

  The number of unused blocks in the `MyISAM` key cache. You can use this value to determine how much of the key cache is in use; see the discussion of `key_buffer_size` in Section 5.1.5, “Server System Variables”.

- `Key_blocks_used`

  The number of used blocks in the `MyISAM` key cache. This value is a high-water mark that indicates the maximum number of blocks that have ever been in use at one time.

- `Key_read_requests`

  The number of requests to read a key block from the `MyISAM` key cache.

- `Key_reads`

  The number of physical reads of a key block from disk into the `MyISAM` key cache. If `Key_reads` is large, then your `key_buffer_size` value is probably too small. The cache miss rate can be calculated as `Key_reads`/`Key_read_requests`.

- `Key_write_requests`

  The number of requests to write a key block to the `MyISAM` key cache.

- `Key_writes`

  The number of physical writes of a key block from the `MyISAM` key cache to disk.

- `Last_query_cost`

  The total cost of the last compiled query as computed by the query optimizer. This is useful for comparing the cost of different query plans for the same query. The default value of 0 means that no query has been compiled yet. The default value is 0. `Last_query_cost` has session scope.

  The `Last_query_cost` value can be computed accurately only for simple “flat” queries, not complex queries such as those with subqueries or `UNION`. For the latter, the value is set to 0.

- `Last_query_partial_plans`

  The number of iterations the query optimizer made in execution plan construction for the previous query. `Last_query_cost` has session scope.

- `Locked_connects`

  The number of attempts to connect to locked user accounts. For information about account locking and unlocking, see Section 6.3.12, “User Account Locking”.

- `Max_execution_time_exceeded`

  The number of `SELECT` statements for which the execution timeout was exceeded.

- `Max_execution_time_set`

  The number of `SELECT` statements for which a nonzero execution timeout was set. This includes statements that include a nonzero `MAX_EXECUTION_TIME` optimizer hint, and statements that include no such hint but execute while the timeout indicated by the `max_execution_time` system variable is nonzero.

- `Max_execution_time_set_failed`

  The number of `SELECT` statements for which the attempt to set an execution timeout failed.

- `Max_used_connections`

  The maximum number of connections that have been in use simultaneously since the server started.

- `Max_used_connections_time`

  The time at which `Max_used_connections` reached its current value.

- `Not_flushed_delayed_rows`

  This status variable is deprecated (because `DELAYED` inserts are not supported), and will be removed in a future release.

- `mecab_charset`

  The character set currently used by the MeCab full-text parser plugin. For related information, see Section 12.9.9, “MeCab Full-Text Parser Plugin”.

- `Ongoing_anonymous_transaction_count`

  Shows the number of ongoing transactions which have been marked as anonymous. This can be used to ensure that no further transactions are waiting to be processed.

- `Ongoing_anonymous_gtid_violating_transaction_count`

  This status variable is only available in debug builds. Shows the number of ongoing transactions which use `gtid_next=ANONYMOUS` and that violate GTID consistency.

- `Ongoing_automatic_gtid_violating_transaction_count`

  This status variable is only available in debug builds. Shows the number of ongoing transactions which use `gtid_next=AUTOMATIC` and that violate GTID consistency.

- `Open_files`

  The number of files that are open. This count includes regular files opened by the server. It does not include other types of files such as sockets or pipes. Also, the count does not include files that storage engines open using their own internal functions rather than asking the server level to do so.

- `Open_streams`

  The number of streams that are open (used mainly for logging).

- `Open_table_definitions`

  The number of cached table definitions.

- `Open_tables`

  The number of tables that are open.

- `Opened_files`

The number of files that have been opened with `my_open()` (a `mysys` library function). Parts of the server that open files without using this function do not increment the count.

- `Opened_table_definitions`

  The number of table definitions that have been cached.

- `Opened_tables`

  The number of tables that have been opened. If `Opened_tables` is big, your `table_open_cache` value is probably too small.

- `Performance_schema_xxx`

  Performance Schema status variables are listed in Section 25.15, “Performance Schema Status Variables”. These variables provide information about instrumentation that could not be loaded or created due to memory constraints.

- `Prepared_stmt_count`

  The current number of prepared statements. (The maximum number of statements is given by the `max_prepared_stmt_count` system variable.)

- `Qcache_free_blocks`

  The number of free memory blocks in the query cache.

- `Qcache_free_memory`

  The amount of free memory for the query cache.

- `Qcache_hits`

  The number of query cache hits.

  The discussion at the beginning of this section indicates how to relate this statement-counting status variable to other such variables.

- `Qcache_inserts`

  The number of queries added to the query cache.

- `Qcache_lowmem_prunes`

  The number of queries that were deleted from the query cache because of low memory.

- `Qcache_not_cached`

  The number of noncached queries (not cacheable, or not cached due to the `query_cache_type` setting).

- `Qcache_queries_in_cache`

  The number of queries registered in the query cache.

- `Qcache_total_blocks`

  The total number of blocks in the query cache.

- `Queries`

The number of statements executed by the server. This variable includes statements executed within stored programs, unlike the `Questions` variable. It does not count `COM_PING` or `COM_STATISTICS` commands.

The discussion at the beginning of this section indicates how to relate this statement-counting status variable to other such variables.

- `Questions`

  The number of statements executed by the server. This includes only statements sent to the server by clients and not statements executed within stored programs, unlike the `Queries` variable. This variable does not count `COM_PING`, `COM_STATISTICS`, `COM_STMT_PREPARE`, `COM_STMT_CLOSE`, or `COM_STMT_RESET` commands.

  The discussion at the beginning of this section indicates how to relate this statement-counting status variable to other such variables.

- `Rpl_semi_sync_master_clients`

  The number of semisynchronous slaves.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_net_avg_wait_time`

  The average time in microseconds the master waited for a slave reply. This variable is always `0`, is deprecated and it will be removed in a future version.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_net_wait_time`

  The total time in microseconds the master waited for slave replies. This variable is always `0`, is deprecated and it will be removed in a future version.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_net_waits`

  The total number of times the master waited for slave replies.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_no_times`

  The number of times the master turned off semisynchronous replication.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_no_tx`

  The number of commits that were not acknowledged successfully by a slave.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_status`

  Whether semisynchronous replication currently is operational on the master. The value is `ON` if the plugin has been enabled and a commit acknowledgment has occurred. It is `OFF` if the plugin is not enabled or the master has fallen back to asynchronous replication due to commit acknowledgment timeout.

This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_timefunc_failures`

  The number of times the master failed when calling time functions such as `gettimeofday()`.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_tx_avg_wait_time`

  The average time in microseconds the master waited for each transaction.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_tx_wait_time`

  The total time in microseconds the master waited for transactions.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_tx_waits`

  The total number of times the master waited for transactions.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_wait_pos_backtraverse`

  The total number of times the master waited for an event with binary coordinates lower than events waited for previously. This can occur when the order in which transactions start waiting for a reply is different from the order in which their binary log events are written.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_wait_sessions`

  The number of sessions currently waiting for slave replies.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_master_yes_tx`

  The number of commits that were acknowledged successfully by a slave.

  This variable is available only if the master-side semisynchronous replication plugin is installed.

- `Rpl_semi_sync_slave_status`

  Whether semisynchronous replication currently is operational on the slave. This is `ON` if the plugin has been enabled and the slave I/O thread is running, `OFF` otherwise.

  This variable is available only if the slave-side semisynchronous replication plugin is installed.

- `Rsa_public_key`

  This variable is available if MySQL was using OpenSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). Its value is the RSA public key value used by the `sha256_password` authentication plugin. The value is nonempty only if the server successfully initializes the private and public keys in the files named by the `sha256_password_private_key_path` and `sha256_password_public_key_path` system variables. The value of `Rsa_public_key` comes from the latter file.

For information about `sha256_password`, see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

- `Select_full_join`

  The number of joins that perform table scans because they do not use indexes. If this value is not 0, you should carefully check the indexes of your tables.

- `Select_full_range_join`

  The number of joins that used a range search on a reference table.

- `Select_range`

  The number of joins that used ranges on the first table. This is normally not a critical issue even if the value is quite large.

- `Select_range_check`

  The number of joins without keys that check for key usage after each row. If this is not 0, you should carefully check the indexes of your tables.

- `Select_scan`

  The number of joins that did a full scan of the first table.

- `Slave_heartbeat_period`

  This variable is obsolete and was removed in MySQL 8.0.1. Instead, use the `HEARTBEAT_INTERVAL` column of the `replication_connection_configuration` table.

- `Slave_last_heartbeat`

  This variable is obsolete and was removed in MySQL 8.0.1. Instead, use the `LAST_HEARTBEAT_TIMESTAMP` column of the `replication_connection_status` table.

- `Slave_open_temp_tables`

  The number of temporary tables that the slave SQL thread currently has open. If the value is greater than zero, it is not safe to shut down the slave; see Section 18.4.1.24, “Replication and Temporary Tables”. This variable reports the total count of open temporary tables for *all* replication channels.

- `Slave_received_heartbeats`

  This variable is obsolete and was removed in MySQL 8.0.1. Instead, use the `COUNT_RECEIVED_HEARTBEATS` column of the `replication_connection_status` table.

- `Slave_retried_transactions`

  This variable is obsolete and was removed in MySQL 8.0.1. Instead, use the `COUNT_TRANSACTIONS_RETRIES` column of the `replication_applier_status` table.

- `Slave_running`

  This variable is obsolete and was removed in MySQL 8.0.1. Instead, use the `SERVICE_STATE` column of the `replication_connection_status` and `replication_applier_status` tables.

- `Slow_launch_threads`

  The number of threads that have taken more than `slow_launch_time` seconds to create.

- `Slow_queries`

The number of queries that have taken more than `long_query_time` seconds. This counter increments regardless of whether the slow query log is enabled. For information about that log, see Section 5.4.5, “The Slow Query Log”.

- `Sort_merge_passes`

  The number of merge passes that the sort algorithm has had to do. If this value is large, you should consider increasing the value of the `sort_buffer_size` system variable.

- `Sort_range`

  The number of sorts that were done using ranges.

- `Sort_rows`

  The number of sorted rows.

- `Sort_scan`

  The number of sorts that were done by scanning the table.

- `Ssl_accept_renegotiates`

  The number of negotiates needed to establish the connection.

- `Ssl_accepts`

  The number of accepted SSL connections.

- `Ssl_callback_cache_hits`

  The number of callback cache hits.

- `Ssl_cipher`

  The current encryption cipher (empty for unencrypted connections).

- `Ssl_cipher_list`

  The list of possible SSL ciphers (empty for non-SSL connections).

- `Ssl_client_connects`

  The number of SSL connection attempts to an SSL-enabled master.

- `Ssl_connect_renegotiates`

  The number of negotiates needed to establish the connection to an SSL-enabled master.

- `Ssl_ctx_verify_depth`

  The SSL context verification depth (how many certificates in the chain are tested).

- `Ssl_ctx_verify_mode`

  The SSL context verification mode.

- `Ssl_default_timeout`

  The default SSL timeout.

- `Ssl_finished_accepts`

  The number of successful SSL connections to the server.

- `Ssl_finished_connects`

  The number of successful slave connections to an SSL-enabled master.

- `Ssl_server_not_after`

  The last date for which the SSL certificate is valid. To check SSL certificate expiration information, use this statement:

```
mysql> SHOW STATUS LIKE 'Ssl_server_not%';
+-----------------------+--------------------------+
| Variable_name         | Value                    |
+-----------------------+--------------------------+
| Ssl_server_not_after  | Apr 28 14:16:39 2025 GMT |
| Ssl_server_not_before | May  1 14:16:39 2015 GMT |
+-----------------------+--------------------------+
```

- `Ssl_server_not_before`

  The first date for which the SSL certificate is valid.

- `Ssl_session_cache_hits`

  The number of SSL session cache hits.

- `Ssl_session_cache_misses`

  The number of SSL session cache misses.

- `Ssl_session_cache_mode`

  The SSL session cache mode.

- `Ssl_session_cache_overflows`

  The number of SSL session cache overflows.

- `Ssl_session_cache_size`

  The SSL session cache size.

- `Ssl_session_cache_timeouts`

  The number of SSL session cache timeouts.

- `Ssl_sessions_reused`

  How many SSL connections were reused from the cache.

- `Ssl_used_session_cache_entries`

  How many SSL session cache entries were used.

- `Ssl_verify_depth`

  The verification depth for replication SSL connections.

- `Ssl_verify_mode`

  The verification mode for replication SSL connections.

- `Ssl_version`

  The SSL protocol version of the connection; for example, TLSv1. If the connection is not encrypted, the value is empty.

- `Table_locks_immediate`

  The number of times that a request for a table lock could be granted immediately.

- `Table_locks_waited`

  The number of times that a request for a table lock could not be granted immediately and a wait was needed. If this is high and you have performance problems, you should first optimize your queries, and then either split your table or tables or use replication.

- `Table_open_cache_hits`

  The number of hits for open tables cache lookups.

- `Table_open_cache_misses`

  The number of misses for open tables cache lookups.

- `Table_open_cache_overflows`

  The number of overflows for the open tables cache. This is the number of times, after a table is opened or closed, a cache instance has an unused entry and the size of the instance is larger than `table_open_cache` / `table_open_cache_instances`.

- `Tc_log_max_pages_used`

  For the memory-mapped implementation of the log that is used by `mysqld` when it acts as the transaction coordinator for recovery of internal XA transactions, this variable indicates the largest number of pages used for the log since the server started. If the product of `Tc_log_max_pages_used` and `Tc_log_page_size` is always significantly less than the log size, the size is larger than necessary and can be reduced. (The size is set by the `--log-tc-size` option. This variable is unused: It is unneeded for binary log-based recovery, and the memory-mapped recovery log method is not used unless the number of storage engines that are capable of two-phase commit and that support XA transactions is greater than one. (`InnoDB` is the only applicable engine.)

- `Tc_log_page_size`

  The page size used for the memory-mapped implementation of the XA recovery log. The default value is determined using `getpagesize()`. This variable is unused for the same reasons as described for `Tc_log_max_pages_used`.

- `Tc_log_page_waits`

  For the memory-mapped implementation of the recovery log, this variable increments each time the server was not able to commit a transaction and had to wait for a free page in the log. If this value is large, you might want to increase the log size (with the `--log-tc-size` option). For binary log-based recovery, this variable increments each time the binary log cannot be closed because there are two-phase commits in progress. (The close operation waits until all such transactions are finished.)

- `Threads_cached`

  The number of threads in the thread cache.

- `Threads_connected`

  The number of currently open connections.

- `Threads_created`

The number of threads created to handle connections. If `Threads_created` is big, you may want to increase the `thread_cache_size` value. The cache miss rate can be calculated as `Threads_created`/`Connections`.

- `Threads_running`

  The number of threads that are not sleeping.

- `Uptime`

  The number of seconds that the server has been up.

- `Uptime_since_flush_status`

  The number of seconds since the most recent `FLUSH STATUS` statement.

5.1.8 Server SQL Modes

The MySQL server can operate in different SQL modes, and can apply these modes differently for different clients, depending on the value of the `sql_mode` system variable. DBAs can set the global SQL mode to match site server operating requirements, and each application can set its session SQL mode to its own requirements.

Modes affect the SQL syntax MySQL supports and the data validation checks it performs. This makes it easier to use MySQL in different environments and to use MySQL together with other database servers.

- Setting the SQL Mode
- The Most Important SQL Modes
- Full List of SQL Modes
- Combination SQL Modes
- Strict SQL Mode
- Comparison of the IGNORE Keyword and Strict SQL Mode
- SQL Mode Changes in MySQL 5.7

For answers to questions often asked about server SQL modes in MySQL, see Section A.3, “MySQL 8.0 FAQ: Server SQL Mode”.

When working with `InnoDB` tables, consider also the `innodb_strict_mode` system variable. It enables additional error checks for `InnoDB` tables.

Setting the SQL Mode

The default SQL mode in MySQL 8.0 includes these modes: `ONLY_FULL_GROUP_BY`, `STRICT_TRANS_TABLES`, `NO_ZERO_IN_DATE`, `NO_ZERO_DATE`, `ERROR_FOR_DIVISION_BY_ZERO`, `NO_AUTO_CREATE_USER`, and `NO_ENGINE_SUBSTITUTION`.

The `ONLY_FULL_GROUP_BY` and `STRICT_TRANS_TABLES` modes were added in MySQL 5.7.5. The `NO_AUTO_CREATE_USER` mode was added in MySQL 5.7.7. The `ERROR_FOR_DIVISION_BY_ZERO`, `NO_ZERO_DATE`, and `NO_ZERO_IN_DATE` modes were added in MySQL 5.7.8. For additional discussion regarding these changes to the default SQL mode value, see SQL Mode Changes in MySQL 5.7.

To set the SQL mode at server startup, use the `--sql-mode="`*`modes`*`"` option on the command line, or `sql-mode="`*`modes`*`"` in an option file such as `my.cnf` (Unix operating systems) or `my.ini` (Windows). *`modes`* is a list of different modes separated by commas. To clear the SQL mode explicitly,

set it to an empty string using `--sql-mode=""` on the command line, or `sql-mode=""` in an option file.

Note

MySQL installation programs may configure the SQL mode during the installation process.

If the SQL mode differs from the default or from what you expect, check for a setting in an option file that the server reads at startup.

To change the SQL mode at runtime, set the global or session `sql_mode` system variable using a `SET` statement:

```
SET GLOBAL sql_mode = 'modes';
SET SESSION sql_mode = 'modes';
```

Setting the `GLOBAL` variable requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege and affects the operation of all clients that connect from that time on. Setting the `SESSION` variable affects only the current client. Each client can change its session `sql_mode` value at any time.

To determine the current global or session `sql_mode` value, use the following statements:

```
SELECT @@GLOBAL.sql_mode;
SELECT @@SESSION.sql_mode;
```

Important

SQL mode and user-defined partitioning. Changing the server SQL mode after creating and inserting data into partitioned tables can cause major changes in the behavior of such tables, and could lead to loss or corruption of data. It is strongly recommended that you never change the SQL mode once you have created tables employing user-defined partitioning.

When replicating partitioned tables, differing SQL modes on master and slave can also lead to problems. For best results, you should always use the same server SQL mode on the master and on the slave.

See Section 22.6, “Restrictions and Limitations on Partitioning”, for more information.

The Most Important SQL Modes

The most important `sql_mode` values are probably these:

- `ANSI`

  This mode changes syntax and behavior to conform more closely to standard SQL. It is one of the special combination modes listed at the end of this section.

- `STRICT_TRANS_TABLES`

  If a value could not be inserted as given into a transactional table, abort the statement. For a nontransactional table, abort the statement if the value occurs in a single-row statement or the first row of a multiple-row statement. More details are given later in this section.

- `TRADITIONAL`

  Make MySQL behave like a “traditional” SQL database system. A simple description of this mode is “give an error instead of a warning” when inserting an incorrect value into a column. It is one of the special combination modes listed at the end of this section.

Note

The `INSERT` or `UPDATE` aborts as soon as the error is noticed. This may not be what you want if you are using a nontransactional storage engine, because data changes made prior to the error may not be rolled back, resulting in a “partially done” update.

When this manual refers to “strict mode,” it means a mode with either or both `STRICT_TRANS_TABLES` or `STRICT_ALL_TABLES` enabled.

Full List of SQL Modes

The following list describes all supported SQL modes:

- `ALLOW_INVALID_DATES`

  Do not perform full checking of dates. Check only that the month is in the range from 1 to 12 and the day is in the range from 1 to 31. This is very convenient for Web applications where you obtain year, month, and day in three different fields and you want to store exactly what the user inserted (without date validation). This mode applies to `DATE` and `DATETIME` columns. It does not apply `TIMESTAMP` columns, which always require a valid date.

  The server requires that month and day values be legal, and not merely in the range 1 to 12 and 1 to 31, respectively. With strict mode disabled, invalid dates such as `'2004-04-31'` are converted to `'0000-00-00'` and a warning is generated. With strict mode enabled, invalid dates generate an error. To permit such dates, enable `ALLOW_INVALID_DATES`.

- `ANSI_QUOTES`

  Treat `"` as an identifier quote character (like the `` ` `` quote character) and not as a string quote character. You can still use `` ` `` to quote identifiers with this mode enabled. With `ANSI_QUOTES` enabled, you cannot use double quotation marks to quote literal strings, because it is interpreted as an identifier.

- `ERROR_FOR_DIVISION_BY_ZERO`

  The `ERROR_FOR_DIVISION_BY_ZERO` mode affects handling of division by zero, which includes `MOD(N,0)`. For data-change operations (`INSERT`, `UPDATE`), its effect also depends on whether strict SQL mode is enabled.

  - If this mode is not enabled, division by zero inserts `NULL` and produces no warning.
  - If this mode is enabled, division by zero inserts `NULL` and produces a warning.
  - If this mode and strict mode are enabled, division by zero produces an error, unless `IGNORE` is given as well. For `INSERT IGNORE` and `UPDATE IGNORE`, division by zero inserts `NULL` and produces a warning.

  For `SELECT`, division by zero returns `NULL`. Enabling `ERROR_FOR_DIVISION_BY_ZERO` causes a warning to be produced as well, regardless of whether strict mode is enabled.

  `ERROR_FOR_DIVISION_BY_ZERO` has an effect when named explicitly and is not part of strict mode. However, it should be used in conjunction with strict mode and is enabled by default. A warning occurs if `ERROR_FOR_DIVISION_BY_ZERO` is enabled without also enabling strict mode or vice versa. For additional discussion, see SQL Mode Changes in MySQL 5.7.

  Because `ERROR_FOR_DIVISION_BY_ZERO` is deprecated, it will be removed in a future MySQL release as a separate mode name and its effect included in the effects of strict SQL mode.

- `HIGH_NOT_PRECEDENCE`

The precedence of the NOT operator is such that expressions such as NOT a BETWEEN b AND c are parsed as NOT (a BETWEEN b AND c). In some older versions of MySQL, the expression was parsed as (NOT a) BETWEEN b AND c. The old higher-precedence behavior can be obtained by enabling the HIGH_NOT_PRECEDENCE SQL mode.

```
mysql> SET sql_mode = '';
mysql> SELECT NOT 1 BETWEEN -5 AND 5;
        -> 0
mysql> SET sql_mode = 'HIGH_NOT_PRECEDENCE';
mysql> SELECT NOT 1 BETWEEN -5 AND 5;
        -> 1
```

- IGNORE_SPACE

  Permit spaces between a function name and the (character. This causes built-in function names to be treated as reserved words. As a result, identifiers that are the same as function names must be quoted as described in Section 9.2, “Schema Object Names”. For example, because there is a COUNT() function, the use of count as a table name in the following statement causes an error:

  ```
  mysql> CREATE TABLE count (i INT);
  ERROR 1064 (42000): You have an error in your SQL syntax
  ```

  The table name should be quoted:

  ```
  mysql> CREATE TABLE `count` (i INT);
  Query OK, 0 rows affected (0.00 sec)
  ```

  The IGNORE_SPACE SQL mode applies to built-in functions, not to user-defined functions or stored functions. It is always permissible to have spaces after a UDF or stored function name, regardless of whether IGNORE_SPACE is enabled.

  For further discussion of IGNORE_SPACE, see Section 9.2.4, “Function Name Parsing and Resolution”.

- NO_AUTO_CREATE_USER

  Prevent the GRANT statement from automatically creating new user accounts if it would otherwise do so, unless authentication information is specified. The statement must specify a nonempty password using IDENTIFIED BY or an authentication plugin using IDENTIFIED WITH.

  It is preferable to create MySQL accounts with CREATE USER rather than GRANT. As of MySQL 5.7.6, NO_AUTO_CREATE_USER is deprecated. As of 5.7.7 the default SQL mode includes NO_AUTO_CREATE_USER and assignments to sql_mode that change the NO_AUTO_CREATE_USER mode state produce a warning, except assignments that set sql_mode to DEFAULT. NO_AUTO_CREATE_USER will be removed in a future MySQL release, at which point its effect will be enabled at all times (GRANT will not create accounts).

- NO_AUTO_VALUE_ON_ZERO

  NO_AUTO_VALUE_ON_ZERO affects handling of AUTO_INCREMENT columns. Normally, you generate the next sequence number for the column by inserting either NULL or 0 into it. NO_AUTO_VALUE_ON_ZERO suppresses this behavior for 0 so that only NULL generates the next sequence number.

  This mode can be useful if 0 has been stored in a table's AUTO_INCREMENT column. (Storing 0 is not a recommended practice, by the way.) For example, if you dump the table with mysqldump and then reload it, MySQL normally generates new sequence numbers when it encounters the 0 values, resulting in a table with contents different from the one that was dumped. Enabling NO_AUTO_VALUE_ON_ZERO before reloading the dump file solves this problem. mysqldump now

automatically includes in its output a statement that enables NO_AUTO_VALUE_ON_ZERO, to avoid this problem.

- NO_BACKSLASH_ESCAPES

  Disable the use of the backslash character (\) as an escape character within strings. With this mode enabled, backslash becomes an ordinary character like any other.

- NO_DIR_IN_CREATE

  When creating a table, ignore all INDEX DIRECTORY and DATA DIRECTORY directives. This option is useful on slave replication servers.

- NO_ENGINE_SUBSTITUTION

  Control automatic substitution of the default storage engine when a statement such as CREATE TABLE or ALTER TABLE specifies a storage engine that is disabled or not compiled in.

  The default SQL mode includes NO_ENGINE_SUBSTITUTION.

  Because storage engines can be pluggable at runtime, unavailable engines are treated the same way:

  With NO_ENGINE_SUBSTITUTION disabled, for CREATE TABLE the default engine is used and a warning occurs if the desired engine is unavailable. For ALTER TABLE, a warning occurs and the table is not altered.

  With NO_ENGINE_SUBSTITUTION enabled, an error occurs and the table is not created or altered if the desired engine is unavailable.

- NO_FIELD_OPTIONS

  Do not print MySQL-specific column options in the output of SHOW CREATE TABLE. This mode is used by mysqldump in portability mode.

- NO_KEY_OPTIONS

  Do not print MySQL-specific index options in the output of SHOW CREATE TABLE. This mode is used by mysqldump in portability mode.

- NO_TABLE_OPTIONS

  Do not print MySQL-specific table options (such as ENGINE) in the output of SHOW CREATE TABLE. This mode is used by mysqldump in portability mode.

- NO_UNSIGNED_SUBTRACTION

  Subtraction between integer values, where one is of type UNSIGNED, produces an unsigned result by default. If the result would otherwise have been negative, an error results:

```
mysql> SET sql_mode = '';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT CAST(0 AS UNSIGNED) - 1;
ERROR 1690 (22003): BIGINT UNSIGNED value is out of range in '(cast(0 as unsigned) - 1)'
```

  If the NO_UNSIGNED_SUBTRACTION SQL mode is enabled, the result is negative:

```
mysql> SET sql_mode = 'NO_UNSIGNED_SUBTRACTION';
mysql> SELECT CAST(0 AS UNSIGNED) - 1;
+--------------------------+
| CAST(0 AS UNSIGNED) - 1 |
+--------------------------+
```

```
|                       -1 |
+--------------------------+
```

If the result of such an operation is used to update an `UNSIGNED` integer column, the result is clipped to the maximum value for the column type, or clipped to 0 if `NO_UNSIGNED_SUBTRACTION` is enabled. If strict SQL mode is enabled, an error occurs and the column remains unchanged.

When `NO_UNSIGNED_SUBTRACTION` is enabled, the subtraction result is signed, *even if any operand is unsigned*. For example, compare the type of column `c2` in table `t1` with that of column `c2` in table `t2`:

```
mysql> SET sql_mode='';
mysql> CREATE TABLE test (c1 BIGINT UNSIGNED NOT NULL);
mysql> CREATE TABLE t1 SELECT c1 - 1 AS c2 FROM test;
mysql> DESCRIBE t1;
+-------+---------------------+------+-----+---------+-------+
| Field | Type                | Null | Key | Default | Extra |
+-------+---------------------+------+-----+---------+-------+
| c2    | bigint(21) unsigned | NO   |     | 0       |       |
+-------+---------------------+------+-----+---------+-------+

mysql> SET sql_mode='NO_UNSIGNED_SUBTRACTION';
mysql> CREATE TABLE t2 SELECT c1 - 1 AS c2 FROM test;
mysql> DESCRIBE t2;
+-------+------------+------+-----+---------+-------+
| Field | Type       | Null | Key | Default | Extra |
+-------+------------+------+-----+---------+-------+
| c2    | bigint(21) | NO   |     | 0       |       |
+-------+------------+------+-----+---------+-------+
```

This means that `BIGINT UNSIGNED` is not 100% usable in all contexts. See Section 12.10, “Cast Functions and Operators”.

- `NO_ZERO_DATE`

  The `NO_ZERO_DATE` mode affects whether the server permits `'0000-00-00'` as a valid date. Its effect also depends on whether strict SQL mode is enabled.

  - If this mode is not enabled, `'0000-00-00'` is permitted and inserts produce no warning.
  - If this mode is enabled, `'0000-00-00'` is permitted and inserts produce a warning.
  - If this mode and strict mode are enabled, `'0000-00-00'` is not permitted and inserts produce an error, unless `IGNORE` is given as well. For `INSERT IGNORE` and `UPDATE IGNORE`, `'0000-00-00'` is permitted and inserts produce a warning.

  As of MySQL 5.7.4, `NO_ZERO_DATE` is deprecated. In MySQL 5.7.4 through 5.7.7, `NO_ZERO_DATE` does nothing when named explicitly. Instead, its effect is included in the effects of strict SQL mode. In MySQL 5.7.8 and later, `NO_ZERO_DATE` does have an effect when named explicitly and is not part of strict mode, as before MySQL 5.7.4. However, it should be used in conjunction with strict mode and is enabled by default. A warning occurs if `NO_ZERO_DATE` is enabled without also enabling strict mode or vice versa. For additional discussion, see SQL Mode Changes in MySQL 5.7.

  Because `NO_ZERO_DATE` is deprecated, it will be removed in a future MySQL release as a separate mode name and its effect included in the effects of strict SQL mode.

- `NO_ZERO_IN_DATE`

  The `NO_ZERO_IN_DATE` mode affects whether the server permits dates in which the year part is nonzero but the month or day part is 0. (This mode affects dates such as `'2010-00-01'` or `'2010-01-00'`, but not `'0000-00-00'`. To control whether the server permits `'0000-00-00'`,

use the NO_ZERO_DATE mode.) The effect of NO_ZERO_IN_DATE also depends on whether strict SQL mode is enabled.

- If this mode is not enabled, dates with zero parts are permitted and inserts produce no warning.
- If this mode is enabled, dates with zero parts are inserted as '0000-00-00' and produce a warning.
- If this mode and strict mode are enabled, dates with zero parts are not permitted and inserts produce an error, unless IGNORE is given as well. For INSERT IGNORE and UPDATE IGNORE, dates with zero parts are inserted as '0000-00-00' and produce a warning.

As of MySQL 5.7.4, NO_ZERO_IN_DATE is deprecated. In MySQL 5.7.4 through 5.7.7, NO_ZERO_IN_DATE does nothing when named explicitly. Instead, its effect is included in the effects of strict SQL mode. In MySQL 5.7.8 and later, NO_ZERO_IN_DATE does have an effect when named explicitly and is not part of strict mode, as before MySQL 5.7.4. However, it should be used in conjunction with strict mode and is enabled by default. A warning occurs if NO_ZERO_IN_DATE is enabled without also enabling strict mode or vice versa. For additional discussion, see SQL Mode Changes in MySQL 5.7.

Because NO_ZERO_IN_DATE is deprecated, it will be removed in a future MySQL release as a separate mode name and its effect included in the effects of strict SQL mode.

- ONLY_FULL_GROUP_BY

  Reject queries for which the select list, HAVING condition, or ORDER BY list refer to nonaggregated columns that are neither named in the GROUP BY clause nor are functionally dependent on (uniquely determined by) GROUP BY columns.

  A MySQL extension to standard SQL permits references in the HAVING clause to aliased expressions in the select list. The HAVING clause can refer to aliases regardless of whether ONLY_FULL_GROUP_BY is enabled.

  For additional discussion and examples, see Section 12.18.3, “MySQL Handling of GROUP BY”.

- PAD_CHAR_TO_FULL_LENGTH

  By default, trailing spaces are trimmed from CHAR column values on retrieval. If PAD_CHAR_TO_FULL_LENGTH is enabled, trimming does not occur and retrieved CHAR values are padded to their full length. This mode does not apply to VARCHAR columns, for which trailing spaces are retained on retrieval.

```
mysql> CREATE TABLE t1 (c1 CHAR(10));
Query OK, 0 rows affected (0.37 sec)

mysql> INSERT INTO t1 (c1) VALUES('xy');
Query OK, 1 row affected (0.01 sec)

mysql> SET sql_mode = '';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT c1, CHAR_LENGTH(c1) FROM t1;
+------+----------------+
| c1   | CHAR_LENGTH(c1) |
+------+----------------+
| xy   |              2 |
+------+----------------+
1 row in set (0.00 sec)

mysql> SET sql_mode = 'PAD_CHAR_TO_FULL_LENGTH';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT c1, CHAR_LENGTH(c1) FROM t1;
```

```
+------------+-----------------+
| c1         | CHAR_LENGTH(c1) |
+------------+-----------------+
| xy         |              10 |
+------------+-----------------+
1 row in set (0.00 sec)
```

- `PIPES_AS_CONCAT`

  Treat `||` as a string concatenation operator (same as `CONCAT()`) rather than as a synonym for `OR`.

- `REAL_AS_FLOAT`

  Treat `REAL` as a synonym for `FLOAT`. By default, MySQL treats `REAL` as a synonym for `DOUBLE`.

- `STRICT_ALL_TABLES`

  Enable strict SQL mode for all storage engines. Invalid data values are rejected. For details, see Strict SQL Mode.

- `STRICT_TRANS_TABLES`

  Enable strict SQL mode for transactional storage engines, and when possible for nontransactional storage engines. For details, see Strict SQL Mode.

- `TIME_TRUNCATE_FRACTIONAL`

  Control whether rounding or truncation occurs when inserting a `TIME`, `DATE`, or `TIMESTAMP` value with a fractional seconds part into a column having the same type but fewer fractional digits. The behavior is to use rounding. If this mode is enabled, truncation occurs instead. The following sequence of statements illustrates the difference:

  ```
  CREATE TABLE t (tval TIME(1));
  SET sql_mode='';
  INSERT INTO t (tval) VALUES(1.55);
  SET sql_mode='TIME_TRUNCATE_FRACTIONAL';
  INSERT INTO t (tval) VALUES(2.55);
  ```

  The resulting table contents look like this, where the first value has been subject to rounding and the second to truncation:

  ```
  mysql> SELECT tval FROM t ORDER BY tval;
  +------------+
  | tval       |
  +------------+
  | 00:00:01.6 |
  | 00:00:02.5 |
  +------------+
  ```

Combination SQL Modes

The following special modes are provided as shorthand for combinations of mode values from the preceding list.

- `ANSI`

  Equivalent to `REAL_AS_FLOAT`, `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, and `ONLY_FULL_GROUP_BY`.

  `ANSI` mode also causes the server to return an error for queries where a set function *`S`* with an outer reference *`S(outer_ref)`* cannot be aggregated in the outer query against which the outer reference has been resolved. This is such a query:

  ```
  ```

```
SELECT * FROM t1 WHERE t1.a IN (SELECT MAX(t1.b) FROM t2 WHERE ...);
```

Here, `MAX(t1.b)` cannot aggregated in the outer query because it appears in the `WHERE` clause of that query. Standard SQL requires an error in this situation. If `ANSI` mode is not enabled, the server treats *`S(outer_ref)`* in such queries the same way that it would interpret *`S(const)`*.

See Section 1.8, “MySQL Standards Compliance”.

- `DB2`

  Equivalent to `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, `NO_KEY_OPTIONS`, `NO_TABLE_OPTIONS`, `NO_FIELD_OPTIONS`.

- `MAXDB`

  Equivalent to `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, `NO_KEY_OPTIONS`, `NO_TABLE_OPTIONS`, `NO_FIELD_OPTIONS`, `NO_AUTO_CREATE_USER`.

- `MSSQL`

  Equivalent to `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, `NO_KEY_OPTIONS`, `NO_TABLE_OPTIONS`, `NO_FIELD_OPTIONS`.

- `MYSQL323`

  Equivalent to `MYSQL323`, `HIGH_NOT_PRECEDENCE`. This means `HIGH_NOT_PRECEDENCE` plus some `SHOW CREATE TABLE` behaviors specific to `MYSQL323`:

  - `TIMESTAMP` column display does not include `DEFAULT` or `ON UPDATE` attributes that were introduced in MySQL 4.1.
  - String column display does not include character set and collation attributes that were introduced in MySQL 4.1. For `CHAR` and `VARCHAR` columns, if the collation is binary, `BINARY` is appended to the column type.
  - The `ENGINE=`*`engine_name`* table option displays as `TYPE=`*`engine+name`*.
  - For `MEMORY` tables, the storage engine is displayed as `HEAP`.

- `MYSQL40`

  Equivalent to `MYSQL40`, `HIGH_NOT_PRECEDENCE`. This means `HIGH_NOT_PRECEDENCE` plus some behaviors specific to `MYSQL40`. These are the same as for `MYSQL323`, except that `SHOW CREATE TABLE` does not display `HEAP` as the storage engine for `MEMORY` tables.

- `ORACLE`

  Equivalent to `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, `NO_KEY_OPTIONS`, `NO_TABLE_OPTIONS`, `NO_FIELD_OPTIONS`, `NO_AUTO_CREATE_USER`.

- `POSTGRESQL`

  Equivalent to `PIPES_AS_CONCAT`, `ANSI_QUOTES`, `IGNORE_SPACE`, `NO_KEY_OPTIONS`, `NO_TABLE_OPTIONS`, `NO_FIELD_OPTIONS`.

- `TRADITIONAL`

  Before MySQL 5.7.4, and in MySQL 5.7.8 and later, `TRADITIONAL` is equivalent to `STRICT_TRANS_TABLES`, `STRICT_ALL_TABLES`, `NO_ZERO_IN_DATE`, `NO_ZERO_DATE`, `ERROR_FOR_DIVISION_BY_ZERO`, `NO_AUTO_CREATE_USER`, and `NO_ENGINE_SUBSTITUTION`.

Strict SQL Mode

Strict mode controls how MySQL handles invalid or missing values in data-change statements such as `INSERT` or `UPDATE`. A value can be invalid for several reasons. For example, it might have the wrong data type for the column, or it might be out of range. A value is missing when a new row to be inserted does not contain a value for a non-`NULL` column that has no explicit `DEFAULT` clause in its definition. (For a `NULL` column, `NULL` is inserted if the value is missing.) Strict mode also affects DDL statements such as `CREATE TABLE`.

If strict mode is not in effect, MySQL inserts adjusted values for invalid or missing values and produces warnings (see Section 13.7.5.40, “SHOW WARNINGS Syntax”). In strict mode, you can produce this behavior by using `INSERT IGNORE` or `UPDATE IGNORE`.

For statements such as `SELECT` that do not change data, invalid values generate a warning in strict mode, not an error.

Strict mode produces an error for attempts to create a key that exceeds the maximum key length. When strict mode is not enabled, this results in a warning and truncation of the key to the maximum key length.

Strict mode does not affect whether foreign key constraints are checked. `foreign_key_checks` can be used for that. (See Section 5.1.5, “Server System Variables”.)

Strict SQL mode is in effect if either `STRICT_ALL_TABLES` or `STRICT_TRANS_TABLES` is enabled, although the effects of these modes differ somewhat:

- For transactional tables, an error occurs for invalid or missing values in a data-change statement when either `STRICT_ALL_TABLES` or `STRICT_TRANS_TABLES` is enabled. The statement is aborted and rolled back.
- For nontransactional tables, the behavior is the same for either mode if the bad value occurs in the first row to be inserted or updated: The statement is aborted and the table remains unchanged. If the statement inserts or modifies multiple rows and the bad value occurs in the second or later row, the result depends on which strict mode is enabled:
  - For `STRICT_ALL_TABLES`, MySQL returns an error and ignores the rest of the rows. However, because the earlier rows have been inserted or updated, the result is a partial update. To avoid this, use single-row statements, which can be aborted without changing the table.
  - For `STRICT_TRANS_TABLES`, MySQL converts an invalid value to the closest valid value for the column and inserts the adjusted value. If a value is missing, MySQL inserts the implicit default value for the column data type. In either case, MySQL generates a warning rather than an error and continues processing the statement. Implicit defaults are described in Section 11.7, “Data Type Default Values”.

Strict mode affects handling of division by zero, zero dates, and zeros in dates as follows:

- Strict mode affects handling of division by zero, which includes `MOD(N,0)`:

  For data-change operations (`INSERT`, `UPDATE`):

  - If strict mode is not enabled, division by zero inserts `NULL` and produces no warning.
  - If strict mode is enabled, division by zero produces an error, unless `IGNORE` is given as well. For `INSERT IGNORE` and `UPDATE IGNORE`, division by zero inserts `NULL` and produces a warning.

  For `SELECT`, division by zero returns `NULL`. Enabling strict mode causes a warning to be produced as well.
- Strict mode affects whether the server permits `'0000-00-00'` as a valid date:
  - If strict mode is not enabled, `'0000-00-00'` is permitted and inserts produce no warning.

  - If strict mode is enabled, `'0000-00-00'` is not permitted and inserts produce an error, unless `IGNORE` is given as well. For `INSERT IGNORE` and `UPDATE IGNORE`, `'0000-00-00'` is permitted and inserts produce a warning.

- Strict mode affects whether the server permits dates in which the year part is nonzero but the month or day part is 0 (dates such as `'2010-00-01'` or `'2010-01-00'`):

  - If strict mode is not enabled, dates with zero parts are permitted and inserts produce no warning.

  - If strict mode is enabled, dates with zero parts are not permitted and inserts produce an error, unless `IGNORE` is given as well. For `INSERT IGNORE` and `UPDATE IGNORE`, dates with zero parts are inserted as `'0000-00-00'` (which is considered valid with `IGNORE`) and produce a warning.

For more information about strict mode with respect to `IGNORE`, see Comparison of the IGNORE Keyword and Strict SQL Mode.

Strict mode affects handling of division by zero, zero dates, and zeros in dates in conjunction with the `ERROR_FOR_DIVISION_BY_ZERO`, `NO_ZERO_DATE`, and `NO_ZERO_IN_DATE` modes.

Comparison of the IGNORE Keyword and Strict SQL Mode

This section compares the effect on statement execution of the `IGNORE` keyword (which downgrades errors to warnings) and strict SQL mode (which upgrades warnings to errors). It describes which statements they affect, and which errors they apply to.

The following table presents a summary comparison of statement behavior when the default is to produce an error versus a warning. An example of when the default is to produce an error is inserting a `NULL` into a `NOT NULL` column. An example of when the default is to produce a warning is inserting a value of the wrong data type into a column (such as inserting the string `'abc'` into an integer column).

Operational Mode	When Statement Default is Error	When Statement Default is Warning
Without `IGNORE` or strict SQL mode	Error	Warning
With `IGNORE`	Warning	Warning (same as without `IGNORE` or strict SQL mode)
With strict SQL mode	Error (same as without `IGNORE` or strict SQL mode)	Error
With `IGNORE` and strict SQL mode	Warning	Warning

One conclusion to draw from the table is that when the `IGNORE` keyword and strict SQL mode are both in effect, `IGNORE` takes precedence. This means that, although `IGNORE` and strict SQL mode can be considered to have opposite effects on error handling, they do not cancel when used together.

The Effect of IGNORE on Statement Execution

Several statements in MySQL support an optional `IGNORE` keyword. This keyword causes the server to downgrade certain types of errors and generate warnings instead. For a multiple-row statement, `IGNORE` causes the statement to skip to the next row instead of aborting.

For example, if the table `t` has a primary key column `i`, attempting to insert the same value of `i` into multiple rows normally produces a duplicate-key error:

```
mysql> INSERT INTO t (i) VALUES(1),(1);
ERROR 1062 (23000): Duplicate entry '1' for key 'PRIMARY'
```

With IGNORE, the row containing the duplicate key still is not inserted, but a warning occurs instead of an error:

```
mysql> INSERT IGNORE INTO t (i) VALUES(1),(1);
Query OK, 1 row affected, 1 warning (0.01 sec)
Records: 2  Duplicates: 1  Warnings: 1

mysql> SHOW WARNINGS;
+---------+------+---------------------------------------+
| Level   | Code | Message                               |
+---------+------+---------------------------------------+
| Warning | 1062 | Duplicate entry '1' for key 'PRIMARY' |
+---------+------+---------------------------------------+
1 row in set (0.00 sec)
```

These statements support the IGNORE keyword:

- CREATE TABLE ... SELECT: IGNORE does not apply to the CREATE TABLE or SELECT parts of the statement but to inserts into the table of rows produced by the SELECT. Rows that duplicate an existing row on a unique key value are discarded.
- DELETE: IGNORE causes MySQL to ignore errors during the process of deleting rows.
- INSERT: With IGNORE, rows that duplicate an existing row on a unique key value are discarded. Rows set to values that would cause data conversion errors are set to the closest valid values instead.

  For partitioned tables where no partition matching a given value is found, IGNORE causes the insert operation to fail silently for rows containing the unmatched value.
- LOAD DATA, LOAD XML: With IGNORE, rows that duplicate an existing row on a unique key value are discarded.
- UPDATE: With IGNORE, rows for which duplicate-key conflicts occur on a unique key value are not updated. Rows updated to values that would cause data conversion errors are updated to the closest valid values instead.

The IGNORE keyword applies to the following errors:

```
ER_BAD_NULL_ERROR
ER_DUP_ENTRY
ER_DUP_ENTRY_WITH_KEY_NAME
ER_DUP_KEY
ER_NO_PARTITION_FOR_GIVEN_VALUE
ER_NO_PARTITION_FOR_GIVEN_VALUE_SILENT
ER_NO_REFERENCED_ROW_2
ER_ROW_DOES_NOT_MATCH_GIVEN_PARTITION_SET
ER_ROW_IS_REFERENCED_2
ER_SUBQUERY_NO_1_ROW
ER_VIEW_CHECK_FAILED
```

The Effect of Strict SQL Mode on Statement Execution

The MySQL server can operate in different SQL modes, and can apply these modes differently for different clients, depending on the value of the sql_mode system variable. In “strict” SQL mode, the server upgrades certain warnings to errors.

For example, in non-strict SQL mode, inserting the string 'abc' into an integer column results in conversion of the value to 0 and a warning:

```
mysql> SET sql_mode = '';
Query OK, 0 rows affected (0.00 sec)
```

```
mysql> INSERT INTO t (i) VALUES('abc');
Query OK, 1 row affected, 1 warning (0.01 sec)

mysql> SHOW WARNINGS;
+---------+------+-------------------------------------------------------+
| Level   | Code | Message                                               |
+---------+------+-------------------------------------------------------+
| Warning | 1366 | Incorrect integer value: 'abc' for column 'i' at row 1 |
+---------+------+-------------------------------------------------------+
1 row in set (0.00 sec)
```

In strict SQL mode, the invalid value is rejected with an error:

```
mysql> SET sql_mode = 'STRICT_ALL_TABLES';
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t (i) VALUES('abc');
ERROR 1366 (HY000): Incorrect integer value: 'abc' for column 'i' at row 1
```

For more information about possible settings of the `sql_mode` system variable, see Section 5.1.8, “Server SQL Modes”.

Strict SQL mode applies to the following statements under conditions for which some value might be out of range or an invalid row is inserted into or deleted from a table:

- `ALTER TABLE`
- `CREATE TABLE`
- `CREATE TABLE ... SELECT`
- `DELETE` (both single table and multiple table)
- `INSERT`
- `LOAD DATA`
- `LOAD XML`
- `SELECT SLEEP()`
- `UPDATE` (both single table and multiple table)

Within stored programs, individual statements of the types just listed execute in strict SQL mode if the program was defined while strict mode was in effect.

Strict SQL mode applies to the following errors, represent a class of errors in which an input value is either invalid or missing. A value is invalid if it has the wrong data type for the column or might be out of range. A value is missing if a new row to be inserted does not contain a value for a `NOT NULL` column that has no explicit `DEFAULT` clause in its definition.

```
ER_BAD_NULL_ERROR
ER_CUT_VALUE_GROUP_CONCAT
ER_DATA_TOO_LONG
ER_DATETIME_FUNCTION_OVERFLOW
ER_DIVISION_BY_ZERO
ER_INVALID_ARGUMENT_FOR_LOGARITHM
ER_NO_DEFAULT_FOR_FIELD
ER_NO_DEFAULT_FOR_VIEW_FIELD
ER_TOO_LONG_KEY
ER_TRUNCATED_WRONG_VALUE
ER_TRUNCATED_WRONG_VALUE_FOR_FIELD
ER_WARN_DATA_OUT_OF_RANGE
ER_WARN_NULL_TO_NOTNULL
ER_WARN_TOO_FEW_RECORDS
ER_WRONG_ARGUMENTS
```

```
ER_WRONG_VALUE_FOR_TYPE
WARN_DATA_TRUNCATED
```

SQL Mode Changes in MySQL 5.7

In MySQL 5.7.5, the ONLY_FULL_GROUP_BY SQL mode is enabled by default because GROUP BY processing has become more sophisticated to include detection of functional dependencies. However, if you find that having ONLY_FULL_GROUP_BY enabled causes queries for existing applications to be rejected, either of these actions should restore operation:

- If it is possible to modify an offending query, do so, either so that nonaggregated columns are functionally dependent on GROUP BY columns, or by referring to nonaggregated columns using ANY_VALUE().
- If it is not possible to modify an offending query (for example, if it is generated by a third-party application), set the sql_mode system variable at server startup to not enable ONLY_FULL_GROUP_BY.

As of MySQL 5.7.4, the ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE SQL modes are deprecated. From MySQL 5.7.4 though 5.7.7, these modes do nothing when named explicitly. Instead, their effects are included in the effects of strict SQL mode (STRICT_ALL_TABLES or STRICT_TRANS_TABLES). In other words, strict mode means the same thing in those versions as the pre-5.7.4 meaning of strict mode plus ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE.

The MySQL 5.7.4 change to make strict mode more strict by including ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE caused some problems. For example, in MySQL 5.6 with strict mode but not NO_ZERO_DATE enabled, TIMESTAMP columns can be defined with DEFAULT '0000-00-00 00:00:00'. In MySQL 5.7.4 with the same mode settings, strict mode includes the effect of NO_ZERO_DATE and TIMESTAMP columns cannot be defined with DEFAULT '0000-00-00 00:00:00'. This causes replication of CREATE TABLE statements from 5.6 to 5.7.4 to fail if they contain such TIMESTAMP columns.

The long term plan is still to have the three affected modes be included in strict SQL mode and to remove them as explicit modes in a future MySQL release. But to restore compatibility in MySQL 5.7 with MySQL 5.6 strict mode and to provide additional time for affected applications to be modified, the following changes were made in MySQL 5.7.8:

- ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE have an effect when named explicitly. This reverts a change made in MySQL 5.7.4.
- ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE are not part of strict SQL mode. This reverts a change made in MySQL 5.7.4.
- ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE are included in the default sql_mode value, which as a result includes these modes: ONLY_FULL_GROUP_BY, STRICT_TRANS_TABLES, NO_ZERO_IN_DATE, NO_ZERO_DATE, ERROR_FOR_DIVISION_BY_ZERO, NO_AUTO_CREATE_USER, and NO_ENGINE_SUBSTITUTION.

With the preceding changes, stricter data checking is still enabled by default, but the individual modes can be disabled in environments where it is currently desirable or necessary to do so.

Although in MySQL 5.7.8 and later ERROR_FOR_DIVISION_BY_ZERO, NO_ZERO_DATE, and NO_ZERO_IN_DATE can be used separately from strict mode, it is intended that they be used together. As a reminder, a warning occurs if they are enabled without also enabling strict mode or vice versa.

Important

The following discussion applies only for MySQL versions 5.7.4 through 5.7.7. For upgrades from a version older than MySQL 5.7.4, we recommend upgrading to MySQL 5.7.8 or later, which renders this discussion unnecessary.

The remainder of this section describes the SQL mode settings to use in MySQL 5.7.4 through 5.7.7 to achieve the same statement execution as before 5.7.4, including the cases for `INSERT` and `UPDATE` in which `IGNORE` is given. It also provides guidelines for determining whether applications need modification to behave the same before and after the SQL mode changes.

The following table shows how to control handling of division by zero for versions other than MySQL 5.7.4 through 5.7.7 and for MySQL 5.7.4 through 5.7.7.

Desired Behavior	MySQL 5.7.x Versions Except 5.7.4 Through 5.7.7	MySQL 5.7.4 Through 5.7.7
insert `NULL`, produce no warning	`ERROR_FOR_DIVISION_BY_ZERO` not enabled	strict mode not enabled
insert `NULL`, produce warning	`ERROR_FOR_DIVISION_BY_ZERO`, or `ERROR_FOR_DIVISION_BY_ZERO` + strict mode + `IGNORE`	strict mode + `IGNORE`
error	`ERROR_FOR_DIVISION_BY_ZERO` + strict mode	strict mode

The following table shows how to control whether the server permits `'0000-00-00'` as a valid date for versions other than MySQL 5.7.4 through 5.7.7 and for MySQL 5.7.4 through 5.7.7.

Desired Behavior	MySQL 5.7.x Versions Except 5.7.4 Through 5.7.7	MySQL 5.7.4 Through 5.7.7
insert `'0000-00-00'`, produce no warning	`NO_ZERO_DATE` not enabled	strict mode not enabled
insert `'0000-00-00'`, produce warning	`NO_ZERO_DATE`, or `NO_ZERO_DATE` + strict mode + `IGNORE`	strict mode + `IGNORE`
error	`NO_ZERO_DATE` + strict mode	strict mode

The following table shows how to control whether the server permits dates with zero parts for versions other than MySQL 5.7.4 through 5.7.7 and for MySQL 5.7.4 through 5.7.7.

Desired Behavior	MySQL 5.7.x Versions Except 5.7.4 Through 5.7.7	MySQL 5.7.4 Through 5.7.7
insert date, produce no warning	`NO_ZERO_IN_DATE` not enabled	strict mode not enabled
insert `'0000-00-00'`, produce warning	`NO_ZERO_IN_DATE`, or `NO_ZERO_IN_DATE` + strict mode + `IGNORE`	strict mode + `IGNORE`
error	`NO_ZERO_IN_DATE` + strict mode	strict mode

The following discussion describes the conditions under which a given statement produces the same or different result under the SQL mode changes in MySQL 5.7.4 through 5.7.7. It considers only strict mode (`STRICT_ALL_TABLES` or `STRICT_TRANS_TABLES`) and the three deprecated modes (`ERROR_FOR_DIVISION_BY_ZERO`, `NO_ZERO_DATE`, and `NO_ZERO_IN_DATE`). Other SQL modes such as `ANSI_QUOTES` or `ONLY_FULL_GROUP_BY` are assumed to be held constant before and after an upgrade.

This discussion also describes how to prepare for an upgrade to 5.7.4 through 5.7.7 from a version older than 5.7.4. *Any modifications should be made before upgrading.*

There is no change in behavior between MySQL 5.6 and 5.7 for the following SQL mode settings. A statement that executes under one of these settings needs no modification to produce the same result in 5.6 and 5.7:

- Strict mode and the three deprecated modes are all not enabled.
- Strict mode and the three deprecated modes are all enabled.

A change from warnings in MySQL 5.6 to no warnings in MySQL 5.7 occurs for the following SQL mode settings. The result of statement execution is the same in 5.6 and 5.7, so statements need no modification unless warnings are considered significant:

- Strict mode is not enabled, but either of the deprecated `ERROR_FOR_DIVISION_BY_ZERO` and `NO_ZERO_DATE` modes are enabled.

A behavior change occurs under the following SQL mode settings. A statement that executes under one of these settings must be modified to produce the same result in 5.6 and 5.7:

- Strict mode is not enabled, `NO_ZERO_IN_DATE` is enabled. For this mode setting, expect these differences in statement execution:
  - In 5.6, the server inserts dates with zero parts as `'0000-00-00'` and produces a warning.
  - In 5.7, the server inserts dates with zero parts as is and produces no warning.
- Strict mode is enabled, with some but not all of the three deprecated modes enabled. For this mode setting, expect these differences in statement execution:

  Statements that would be affected by enabling the not-enabled deprecated modes produce errors in 5.7 but not in 5.6. Suppose that strict mode, `NO_ZERO_DATE`, and `NO_ZERO_IN_DATE` are enabled, and a data-change statement performs division by zero:

  - In 5.6, the statement inserts `NULL` and produces no warning. Enabling `ERROR_FOR_DIVISION_BY_ZERO` would cause an error instead.
  - In 5.7, an error occurs because strict mode implicitly includes the effect of `ERROR_FOR_DIVISION_BY_ZERO`. Enabling `ERROR_FOR_DIVISION_BY_ZERO` explicitly would not change that.

To prepare for an upgrade to MySQL 5.7.4 through 5.7.7, the main principle is to make sure that your applications will operate the same way in MySQL 5.6 and 5.7. For example, you can adopt either of these approaches to application compatibility:

- Modify the application to set the SQL mode on a version-specific basis. If we assume that an application will not be used with development versions of MySQL 5.7 prior to 5.7.4, it is possible to set the `sql_mode` value for the application based on the current server version as follows:

  ```
  SET sql_mode = IF(LEFT(VERSION(),3)<'5.7',5.6 mode,5.7 mode);
  ```

  The tables shown earlier in this section serve as a guide to the appropriate equivalent modes for MySQL 5.6 and 5.7.
- Modify the application to execute under an SQL mode for which statements produce the same result in MySQL 5.6 and 5.7.

Tip

`TRADITIONAL` SQL mode in MySQL 5.6 includes strict mode and the three deprecated modes. If you write applications to operate in `TRADITIONAL` mode in MySQL 5.6, there is no change to make for MySQL 5.7.

When assessing SQL mode compatibility between MySQL 5.6 and 5.7, consider particularly these statement execution contexts:

- Replication. You will encounter replication incompatibility related to the SQL mode changes under the following conditions:

- MySQL 5.6 master and 5.7 slave
- Statement-based replication
- An SQL mode setting for which statements produce different results in MySQL 5.6 and 5.7, as described earlier

To handle this incompatibility, use one of these workarounds:

- Use row-based replication
- Use `IGNORE`
- Use an SQL mode for which statements do not produce different results in MySQL 5.6 and 5.7

- Stored programs (stored procedures and functions, triggers, and events). Each stored program executes using the SQL mode in effect at the time it was created. To identify stored programs that may be affected by differences between MySQL 5.6 and 5.7 in SQL mode handling, use these queries:

```
SELECT ROUTINE_SCHEMA, ROUTINE_NAME, ROUTINE_TYPE, SQL_MODE
FROM INFORMATION_SCHEMA.ROUTINES
WHERE SQL_MODE LIKE '%STRICT%'
OR SQL_MODE LIKE '%DIVISION%'
OR SQL_MODE LIKE '%NO_ZERO%';

SELECT TRIGGER_SCHEMA, TRIGGER_NAME, SQL_MODE
FROM INFORMATION_SCHEMA.TRIGGERS
WHERE SQL_MODE LIKE '%STRICT%'
OR SQL_MODE LIKE '%DIVISION%'
OR SQL_MODE LIKE '%NO_ZERO%';

SELECT EVENT_SCHEMA, EVENT_NAME, SQL_MODE
FROM INFORMATION_SCHEMA.EVENTS
WHERE SQL_MODE LIKE '%STRICT%'
OR SQL_MODE LIKE '%DIVISION%'
OR SQL_MODE LIKE '%NO_ZERO%';
```

5.1.9 IPv6 Support

Support for IPv6 in MySQL includes these capabilities:

- MySQL Server can accept TCP/IP connections from clients connecting over IPv6. For example, this command connects over IPv6 to the MySQL server on the local host:

```
shell> mysql -h ::1
```

To use this capability, two things must be true:

- Your system must be configured to support IPv6. See Section 5.1.9.1, “Verifying System Support for IPv6”.
- The default MySQL server configuration permits IPv6 connections in addition to IPv4 connections. To change the default configuration, start the server with an appropriate `--bind-address` option. See Section 5.1.5, “Server System Variables”.

- MySQL account names permit IPv6 addresses to enable DBAs to specify privileges for clients that connect to the server over IPv6. See Section 6.2.4, “Specifying Account Names”. IPv6 addresses can be specified in account names in statements such as `CREATE USER`, `GRANT`, and `REVOKE`. For example:

```
mysql> CREATE USER 'bill'@'::1' IDENTIFIED BY 'secret';
mysql> GRANT SELECT ON mydb.* TO 'bill'@'::1';
```

- IPv6 functions enable conversion between string and internal format IPv6 address formats, and checking whether values represent valid IPv6 addresses. For example, `INET6_ATON()` and `INET6_NTOA()` are similar to `INET_ATON()` and `INET_NTOA()`, but handle IPv6 addresses in addition to IPv4 addresses. See Section 12.20, “Miscellaneous Functions”.

The following sections describe how to set up MySQL so that clients can connect to the server over IPv6.

5.1.9.1 Verifying System Support for IPv6

Before MySQL Server can accept IPv6 connections, the operating system on your server host must support IPv6. As a simple test to determine whether that is true, try this command:

```
shell> ping6 ::1
16 bytes from ::1, icmp_seq=0 hlim=64 time=0.171 ms
16 bytes from ::1, icmp_seq=1 hlim=64 time=0.077 ms
...
```

To produce a description of your system's network interfaces, invoke `ifconfig -a` and look for IPv6 addresses in the output.

If your host does not support IPv6, consult your system documentation for instructions on enabling it. It might be that you need only reconfigure an existing network interface to add an IPv6 address. Or a more extensive change might be needed, such as rebuilding the kernel with IPv6 options enabled.

These links may be helpful in setting up IPv6 on various platforms:

- Windows
- Gentoo Linux
- Ubuntu Linux
- Linux (Generic)
- OS X

5.1.9.2 Configuring the MySQL Server to Permit IPv6 Connections

The MySQL server listens on a single network socket for TCP/IP connections. This socket is bound to a single address, but it is possible for an address to map onto multiple network interfaces. To specify an address, use the `--bind-address=addr` option at server startup, where *`addr`* is an IPv4 or IPv6 address or a host name. If *`addr`* is a host name, the server resolves the name to an IP address and binds to that address.

The server treats different types of addresses as follows:

- If the address is `*`, the server accepts TCP/IP connections on all server host IPv6 and IPv4 interfaces if the server host supports IPv6, or accepts TCP/IP connections on all IPv4 addresses otherwise. Use this address to permit both IPv4 and IPv6 connections on all server interfaces. This value is the default.
- If the address is `0.0.0.0`, the server accepts TCP/IP connections on all server host IPv4 interfaces.
- If the address is `::`, the server accepts TCP/IP connections on all server host IPv4 and IPv6 interfaces. Use this address to permit both IPv4 and IPv6 connections on all server interfaces.
- If the address is an IPv4-mapped address, the server accepts TCP/IP connections for that address, in either IPv4 or IPv6 format. For example, if the server is bound to `::ffff:127.0.0.1`, clients can connect using `--host=127.0.0.1` or `--host=::ffff:127.0.0.1`.

- If the address is a “regular” IPv4 or IPv6 address (such as `127.0.0.1` or `::1`), the server accepts TCP/IP connections only for that IPv4 or IPv6 address.

If you intend to bind the server to a specific address, be sure that the `mysql.user` grant table contains an account with administrative privileges that you can use to connect to that address. Otherwise, you will not be able to shut down the server. For example, if you bind the server to `*`, you can connect to it using all existing accounts. But if you bind the server to `::1`, it accepts connections only on that address. In that case, first make sure that the `'root'@'::1'` account is present in the `mysql.user` table so you can still connect to the server to shut it down.

5.1.9.3 Connecting Using the IPv6 Local Host Address

The following procedure shows how to configure MySQL to permit IPv6 connections by clients that connect to the local server using the `::1` local host address. The instructions given here assume that your system supports IPv6.

1. Start the MySQL server with an appropriate `--bind-address` option to permit it to accept IPv6 connections. For example, put the following lines in your server option file and restart the server:

   ```
   [mysqld]
   bind-address = *
   ```

   Alternatively, you can bind the server to `::1`, but that makes the server more restrictive for TCP/IP connections. It accepts only IPv6 connections for that single address and rejects IPv4 connections. For more information, see Section 5.1.9.2, “Configuring the MySQL Server to Permit IPv6 Connections”.

2. As an administrator, connect to the server and create an account for a local user who will connect from the `::1` local IPv6 host address:

   ```
   mysql> CREATE USER 'ipv6user'@'::1' IDENTIFIED BY 'ipv6pass';
   ```

   For the permitted syntax of IPv6 addresses in account names, see Section 6.2.4, “Specifying Account Names”. In addition to the `CREATE USER` statement, you can issue `GRANT` statements that give specific privileges to the account, although that is not necessary for the remaining steps in this procedure.

3. Invoke the `mysql` client to connect to the server using the new account:

   ```
   shell> mysql -h ::1 -u ipv6user -pipv6pass
   ```

4. Try some simple statements that show connection information:

   ```
   mysql> STATUS
   ...
   Connection:   ::1 via TCP/IP
   ...

   mysql> SELECT CURRENT_USER(), @@bind_address;
   +----------------+----------------+
   | CURRENT_USER() | @@bind_address |
   +----------------+----------------+
   | ipv6user@::1   | ::             |
   +----------------+----------------+
   ```

5.1.9.4 Connecting Using IPv6 Nonlocal Host Addresses

The following procedure shows how to configure MySQL to permit IPv6 connections by remote clients. It is similar to the preceding procedure for local clients, but the server and client hosts are distinct and each has its own nonlocal IPv6 address. The example uses these addresses:

```
Server host: 2001:db8:0:f101::1
Client host: 2001:db8:0:f101::2
```

These addresses are chosen from the nonroutable address range recommended by IANA for documentation purposes and suffice for testing on your local network. To accept IPv6 connections from clients outside the local network, the server host must have a public address. If your network provider assigns you an IPv6 address, you can use that. Otherwise, another way to obtain an address is to use an IPv6 broker; see Section 5.1.9.5, “Obtaining an IPv6 Address from a Broker”.

1. Start the MySQL server with an appropriate `--bind-address` option to permit it to accept IPv6 connections. For example, put the following lines in your server option file and restart the server:

   ```
   [mysqld]
   bind-address = *
   ```

   Alternatively, you can bind the server to `2001:db8:0:f101::1`, but that makes the server more restrictive for TCP/IP connections. It accepts only IPv6 connections for that single address and rejects IPv4 connections. For more information, see Section 5.1.9.2, “Configuring the MySQL Server to Permit IPv6 Connections”.

2. On the server host (`2001:db8:0:f101::1`), create an account for a user who will connect from the client host (`2001:db8:0:f101::2`):

   ```
   mysql> CREATE USER 'remoteipv6user'@'2001:db8:0:f101::2' IDENTIFIED BY 'remoteipv6pass';
   ```

3. On the client host (`2001:db8:0:f101::2`), invoke the `mysql` client to connect to the server using the new account:

   ```
   shell> mysql -h 2001:db8:0:f101::1 -u remoteipv6user -premoteipv6pass
   ```

4. Try some simple statements that show connection information:

   ```
   mysql> STATUS
   ...
   Connection:   2001:db8:0:f101::1 via TCP/IP
   ...

   mysql> SELECT CURRENT_USER(), @@bind_address;
   +----------------------------------+----------------+
   | CURRENT_USER()                   | @@bind_address |
   +----------------------------------+----------------+
   | remoteipv6user@2001:db8:0:f101::2 | ::             |
   +----------------------------------+----------------+
   ```

5.1.9.5 Obtaining an IPv6 Address from a Broker

If you do not have a public IPv6 address that enables your system to communicate over IPv6 outside your local network, you can obtain one from an IPv6 broker. The Wikipedia IPv6 Tunnel Broker page lists several brokers and their features, such as whether they provide static addresses and the supported routing protocols.

After configuring your server host to use a broker-supplied IPv6 address, start the MySQL server with an appropriate `--bind-address` option to permit the server to accept IPv6 connections. For example, put the following lines in the server option file and restart the server:

```
[mysqld]
bind-address = *
```

Alternatively, you can bind the server to the specific IPv6 address provided by the broker, but that makes the server more restrictive for TCP/IP connections. It accepts only IPv6 connections for that

single address and rejects IPv4 connections. For more information, see Section 5.1.9.2, “Configuring the MySQL Server to Permit IPv6 Connections”. In addition, if the broker allocates dynamic addresses, the address provided for your system might change the next time you connect to the broker. If so, any accounts you create that name the original address become invalid. To bind to a specific address but avoid this change-of-address problem, you may be able to arrange with the broker for a static IPv6 address.

The following example shows how to use Freenet6 as the broker and the `gogoc` IPv6 client package on Gentoo Linux.

1. Create an account at Freenet6 by visiting this URL and signing up:

   ```
   http://gogonet.gogo6.com
   ```

2. After creating the account, go to this URL, sign in, and create a user ID and password for the IPv6 broker:

   ```
   http://gogonet.gogo6.com/page/freenet6-registration
   ```

3. As `root`, install `gogoc`:

   ```
   shell> emerge gogoc
   ```

4. Edit `/etc/gogoc/gogoc.conf` to set the `userid` and `password` values. For example:

   ```
   userid=gogouser
   passwd=gogopass
   ```

5. Start `gogoc`:

   ```
   shell> /etc/init.d/gogoc start
   ```

   To start `gogoc` each time your system boots, execute this command:

   ```
   shell> rc-update add gogoc default
   ```

6. Use `ping6` to try to ping a host:

   ```
   shell> ping6 ipv6.google.com
   ```

7. To see your IPv6 address:

   ```
   shell> ifconfig tun
   ```

5.1.10 Server-Side Help

MySQL Server supports a `HELP` statement that returns information from the MySQL Reference manual (see Section 13.8.3, “HELP Syntax”). Several tables in the `mysql` system database contain the information needed to support this statement (see Section 5.3, “The mysql System Database”). The proper operation of this statement requires that these help tables be initialized, which is done by processing the contents of the `fill_help_tables.sql` script.

If you install MySQL using a binary or source distribution on Unix, help table content initialization occurs when you initialize the data directory (see Section 2.9.1, “Initializing the Data Directory”). For an RPM distribution on Linux or binary distribution on Windows, content initialization occurs as part of the MySQL installation process.

If you upgrade MySQL using a binary distribution, help table content is not upgraded automatically, but you can upgrade it manually. Locate the `fill_help_tables.sql` file in the `share` or `share/mysql` directory. Change location into that directory and process the file with the `mysql` client as follows:

```
shell> mysql -u root mysql < fill_help_tables.sql
```

You can also obtain the latest `fill_help_tables.sql` at any time to upgrade your help tables. Download the proper file for your version of MySQL from http://dev.mysql.com/doc/index-other.html. After downloading and uncompressing the file, process it with `mysql` as described previously.

If you are working with Git and a MySQL development source tree, you must use a downloaded copy of the `fill_help_tables.sql` file because the source tree contains only a “stub” version.

Note

For a server that participates in replication, the help table content upgrade process involves multiple servers. For details, see Section 18.4.1.29, “Replication of Server-Side Help Tables”.

5.1.11 Server Response to Signals

On Unix, signals can be sent to processes. `mysqld` responds to signals sent to it as follows:

- `SIGTERM` causes the server to shut down.
- `SIGHUP` causes the server to reload the grant tables and to flush tables, logs, the thread cache, and the host cache. These actions are like various forms of the `FLUSH` statement. The server also writes a status report to the error log that has this format:

```
Status information:

Current dir: /var/mysql/data/
Running threads: 0  Stack size: 196608
Current locks:

Key caches:
default
Buffer_size:       8388600
Block_size:           1024
Division_limit:        100
Age_limit:             300
blocks used:             0
not flushed:             0
w_requests:              0
writes:                  0
r_requests:              0
reads:                   0

handler status:
read_key:            0
read_next:           0
read_rnd             0
read_first:          1
write:               0
delete               0
update:              0

Table status:
Opened tables:          5
Open tables:            0
Open files:             7
Open streams:           0

Alarm status:
Active alarms:   1
Max used alarms: 2
```

```
Next alarm time: 67
```

5.1.12 The Server Shutdown Process

The server shutdown process takes place as follows:

1. The shutdown process is initiated.

   This can occur initiated several ways. For example, a user with the `SHUTDOWN` privilege can execute a `mysqladmin shutdown` command. `mysqladmin` can be used on any platform supported by MySQL. Other operating system-specific shutdown initiation methods are possible as well: The server shuts down on Unix when it receives a `SIGTERM` signal. A server running as a service on Windows shuts down when the services manager tells it to.

2. The server creates a shutdown thread if necessary.

   Depending on how shutdown was initiated, the server might create a thread to handle the shutdown process. If shutdown was requested by a client, a shutdown thread is created. If shutdown is the result of receiving a `SIGTERM` signal, the signal thread might handle shutdown itself, or it might create a separate thread to do so. If the server tries to create a shutdown thread and cannot (for example, if memory is exhausted), it issues a diagnostic message that appears in the error log:

   ```
   Error: Can't create thread to kill server
   ```

3. The server stops accepting new connections.

   To prevent new activity from being initiated during shutdown, the server stops accepting new client connections by closing the handlers for the network interfaces to which it normally listens for connections: the TCP/IP port, the Unix socket file, the Windows named pipe, and shared memory on Windows.

4. The server terminates current activity.

   For each thread associated with a client connection, the server breaks the connection to the client and marks the thread as killed. Threads die when they notice that they are so marked. Threads for idle connections die quickly. Threads that currently are processing statements check their state periodically and take longer to die. For additional information about thread termination, see Section 13.7.6.4, “KILL Syntax”, in particular for the instructions about killed `REPAIR TABLE` or `OPTIMIZE TABLE` operations on `MyISAM` tables.

   For threads that have an open transaction, the transaction is rolled back. If a thread is updating a nontransactional table, an operation such as a multiple-row `UPDATE` or `INSERT` may leave the table partially updated because the operation can terminate before completion.

   If the server is a master replication server, it treats threads associated with currently connected slaves like other client threads. That is, each one is marked as killed and exits when it next checks its state.

   If the server is a slave replication server, it stops the I/O and SQL threads, if they are active, before marking client threads as killed. The SQL thread is permitted to finish its current statement (to avoid causing replication problems), and then stops. If the SQL thread is in the middle of a transaction at this point, the server waits until the current replication event group (if any) has finished executing, or until the user issues a `KILL QUERY` or `KILL CONNECTION` statement. See also Section 13.4.2.7, “STOP SLAVE Syntax”. Since nontransactional statements cannot be rolled back, in order to guarantee crash-safe replication, only transactional tables should be used.

Note

To guarantee crash safety on the slave, you must run the slave with `--relay-log-recovery` enabled.

See also Section 18.2.4, “Replication Relay and Status Logs”).

5. The server shuts down or closes storage engines.

   At this stage, the server flushes the table cache and closes all open tables.

   Each storage engine performs any actions necessary for tables that it manages. `InnoDB` flushes its buffer pool to disk (unless `innodb_fast_shutdown` is 2), writes the current LSN to the tablespace, and terminates its own internal threads. `MyISAM` flushes any pending index writes for a table.

6. The server exits.

To provide information to management processes, the server returns one of the exit codes described in the following list. The phrase in parentheses indicates the action taken by systemd in response to the code, for platforms on which systemd is used to manage the server.

- 0 = successful termination (no restart done)
- 1 = unsuccessful termination (no restart done)
- 2 = unsuccessful termination (restart done)

5.2 The MySQL Data Directory

Information managed by the MySQL server is stored under a directory known as the data directory. The following list briefly describes the items typically found in the data directory, with cross references for additional information:

- Data directory subdirectories. Each subdirectory of the data directory is a database directory and corresponds to a database managed by the server. All MySQL installations have certain standard databases:
  - The `mysql` directory corresponds to the `mysql` system database, which contains information required by the MySQL server as it runs. This database contains data dictionary tables and system tables. See Section 5.3, “The mysql System Database”.
  - The `performance_schema` directory corresponds to the Performance Schema, which provides information used to inspect the internal execution of the server at runtime. See Chapter 25, *MySQL Performance Schema*.
  - The `sys` directory corresponds to the `sys` schema, which provides a set of objects to help interpret Performance Schema information more easily. See Chapter 26, *MySQL sys Schema*.

  Other subdirectories correspond to databases created by users or applications.

  Note

  `INFORMATION_SCHEMA` is a standard database, but its implementation uses no corresponding database directory.

- Log files written by the server. See Section 5.4, “MySQL Server Logs”.
- `InnoDB` tablespace and log files. See Chapter 15, *The InnoDB Storage Engine*.
- Default/autogenerated SSL and RSA certificate and key files. See Section 6.4.6, “Creating SSL and RSA Certificates and Keys”.
- The server process ID file (while the server is running).
- The `mysqld-auto.cnf` file that stores persisted global system variable settings. See Section 5.1.6, “Using System Variables”.

Some items in the preceding list can be relocated elsewhere by reconfiguring the server. In addition, the `--datadir` option enables the location of the data directory itself to be changed. For a given MySQL installation, check the server configuration to determine whether items have been moved.

5.3 The mysql System Database

The `mysql` database is the system database. It contains tables that store information required by the MySQL server as it runs. A broad categorization is that the `mysql` database contains data dictionary tables that store database object metadata, and system tables used for other operational purposes. The following discussion further subdivides the set of system tables into smaller categories.

- Data Dictionary Tables
- Grant System Tables
- Object Information System Tables
- Log System Tables
- Server-Side Help System Tables
- Time Zone System Tables
- Replication System Tables
- Optimizer System Tables
- Miscellaneous System Tables

The remainder of this section enumerates the tables in each category, with cross references for additional information. Data dictionary tables and system tables use the `InnoDB` storage engine unless otherwise indicated.

As of MySQL 8.0.2, `mysql` system tables and data dictionary tables reside in a single `InnoDB` tablespace file named `mysql.ibd` in the MySQL data directory. Previously, these tables were created in individual tablespace files in the `mysql` database directory.

Data Dictionary Tables

These tables comprise the data dictionary, which contains metadata about database objects. For additional information, see Chapter 14, *MySQL Data Dictionary*.

Important

The data dictionary is new in MySQL 8.0. A data dictionary-enabled server entails some general operational differences compared to previous MySQL releases. For details, see Section 14.6, “Data Dictionary Usage Differences”. Also, for upgrades to MySQL 8.0 from MySQL 5.7, the upgrade procedure differs somewhat from previous MySQL releases and requires that you verify the upgrade readiness of your installation by checking specific prerequisites. For more information, see Section 2.10.1, “Upgrading MySQL”, particularly Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation.

- `catalogs`: Catalog information.
- `character_sets`: Information about available character sets.
- `collations`: Information about collations for each character set.
- `column_statistics`: Histogram statistics for column values. See Section 8.9.6, “Optimizer Statistics”.
- `column_type_elements`: Information about types used by columns.

- `columns`: Information about columns in tables.
- `dd_properties`: A table that identifies data dictionary properties, such as its version. The server uses this to determine whether the data dictionary must be upgraded to a newer version.
- `events`: Information about Event Scheduler events. See Section 23.4, “Using the Event Scheduler”. The server loads events listed in this table during its startup sequence, unless started with the `--skip-grant-tables` option.
- `foreign_keys`, `foreign_key_column_usage`: Information about foreign keys.
- `index_column_usage`: Information about columns used by indexes.
- `index_partitions`: Information about partitions used by indexes.
- `index_stats`: Used to store dynamic index statistics generated when `ANALYZE TABLE` is executed.
- `indexes`: Information about table indexes.
- `parameter_type_elements`: Information about stored procedure and function parameters, and about return values for stored functions.
- `parameters`: Information about stored procedures and functions. See Section 23.2, “Using Stored Routines (Procedures and Functions)”.
- `routines`: Information about stored procedures and functions. See Section 23.2, “Using Stored Routines (Procedures and Functions)”.
- `schemata`: Information about schemata. In MySQL, a schema is a database, so this table provides information about databases.
- `st_spatial_reference_systems`: Information about available spatial reference systems for spatial data.
- `table_partition_values`: Information about values used by table partitions.
- `table_partitions`: Information about partitions used by tables.
- `table_stats`: Information about dynamic table statistics generated when `ANALYZE TABLE` is executed.
- `tables`: Information about tables in databases.
- `tablespace_files`: Information about files used by tablespaces.
- `tablespaces`: Information about active tablespaces.
- `triggers`: Information about triggers.
- `version`: A single-row table that identifies the data dictionary version. The server uses this to determine whether the data dictionary must be upgraded to a newer version. In MySQL 8.0.2, this table was replaced by the `dd_properties` data dictionary table.
- `view_routine_usage`: Information about dependencies between views and stored functions used by them.
- `view_table_usage`: Used to track dependencies between views and their underlying tables.

Data dictionary tables are invisible. They cannot be read with `SELECT`, do not appear in the output of `SHOW TABLES`, are not listed in the `INFORMATION_SCHEMA.TABLES` table, and so forth. However, in most cases there are corresponding `INFORMATION_SCHEMA` tables that can be queried. Conceptually, the `INFORMATION_SCHEMA` provides a view through which MySQL exposes data dictionary metadata. For example, you cannot select from the `mysql.schemata` table directly:

```
mysql> SELECT * FROM mysql.schemata;
ERROR 3554 (HY000): Access to data dictionary table 'mysql.schemata' is rejected.
```

Instead, select that information from the corresponding `INFORMATION_SCHEMA` table:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.SCHEMATA\G
*************************** 1. row ***************************
              CATALOG_NAME: def
               SCHEMA_NAME: mysql
DEFAULT_CHARACTER_SET_NAME: utf8mb4
    DEFAULT_COLLATION_NAME: utf8mb4_0900_ai_ci
                  SQL_PATH: NULL
*************************** 2. row ***************************
              CATALOG_NAME: def
               SCHEMA_NAME: information_schema
DEFAULT_CHARACTER_SET_NAME: utf8
    DEFAULT_COLLATION_NAME: utf8_general_ci
                  SQL_PATH: NULL
...
```

There is no `INFORMATION_SCHEMA` table that corresponds exactly to `mysql.indexes`, but `INFORMATION_SCHEMA.STATISTICS` contains much of the same information.

As of yet, there are no `INFORMATION_SCHEMA` tables that correspond exactly to `mysql.foreign_keys`, `mysql.foreign_key_column_usage`. The standard SQL way to obtain foreign key information is by using the `INFORMATION_SCHEMA REFERENTIAL_CONSTRAINTS` and `KEY_COLUMN_USAGE` tables; these tables are now implemented as views on the `foreign_keys`, `foreign_key_column_usage`, and other data dictionary tables.

Some system tables from before MySQL 8.0 have been replaced by data dictionary tables and are no longer present in the `mysql` system database:

- The `events` data dictionary table supersedes the `event` table from before MySQL 8.0.
- The `parameters` and `routines` data dictionary tables together supersede the `proc` table from before MySQL 8.0.

Grant System Tables

These system tables contain grant information about user accounts and the privileges held by them. For additional information about the structure, contents, and purpose of the these tables, see Section 6.2.3, “Grant Tables”.

As of MySQL 8.0, the grant tables are `InnoDB` (transactional) tables. Previously, these were `MyISAM` (nontransactional) tables. The change of grant-table storage engine underlies an accompanying change in MySQL 8.0 to the behavior of account-management statements such as `CREATE USER` and `GRANT`. Previously, an account-management statement that named multiple users could succeed for some users and fail for others. The statements are now transactional and either succeed for all named users or roll back and have no effect if any error occurs.

Note

If MySQL is upgraded from an older version but the grant tables have not been upgraded from `MyISAM` to `InnoDB`, the server considers them read only and account-management statements produce an error. For upgrade instructions, see Section 2.10.1, “Upgrading MySQL”.

- `user`: User accounts, global privileges, and other non-privilege columns.
- `global_grants`: Assignments of dynamic global privileges to users; see Section 6.2.2, “Static Versus Dynamic Privileges”.
- `db`: Database-level privileges.

- `tables_priv`: Table-level privileges.
- `columns_priv`: Column-level privileges.
- `procs_priv`: Stored procedure and function privileges.
- `proxies_priv`: Proxy-user privileges.
- `default_roles`: This table lists default roles to be activated after a user connects and authenticates, or executes `SET ROLE DEFAULT`.
- `role_edges`: This table lists edges for role subgraphs.

  A given `user` table row might refer to a user account or a role. The server can distinquish whether a row represents a user account, a role, or both by consulting the `role_edges` table for information about relations between authentication IDs.

Object Information System Tables

These system tables contain information about stored programs, components, user-defined functions, and server-side plugins:

- `component`: The registry for server components. Any components listed in this table are installed by a loader service during the server startup sequence. See Section 5.5, “MySQL Server Components”.
- `func`: Information about user-defined functions (UDFs). See Section 28.4, “Adding New Functions to MySQL”. The server loads UDFs listed in this table during its startup sequence, unless started with the `--skip-grant-tables` option.
- `plugin`: Information about server-side plugins. See Section 5.6.2, “Installing and Uninstalling Plugins”, and Section 28.2, “The MySQL Plugin API”. The server loads plugins listed in this table during its startup sequence, unless started with the `--skip-grant-tables` option.

Log System Tables

The server uses these system tables for logging:

- `general_log`: The general query log table.
- `slow_log`: The slow query log table.

Log tables use the `CSV` storage engine.

For more information, see Section 5.4, “MySQL Server Logs”.

Server-Side Help System Tables

These system tables contain server-side help information:

- `help_category`: Information about help categories.
- `help_keyword`: Keywords associated with help topics.
- `help_relation`: Mappings between help keywords and topics.
- `help_topic`: Help topic contents.

For more information, see Section 5.1.10, “Server-Side Help”.

Time Zone System Tables

These system tables contain time zone information:

- `time_zone`: Time zone IDs and whether they use leap seconds.

- `time_zone_leap_second`: When leap seconds occur.
- `time_zone_name`: Mappings between time zone IDs and names.
- `time_zone_transition`, `time_zone_transition_type`: Time zone descriptions.

For more information, see Section 10.6, “MySQL Server Time Zone Support”.

Replication System Tables

The server uses these system tables to support replication:

- `gtid_executed`: Table for storing GTID values. See mysql.gtid_executed Table.
- `ndb_binlog_index`: Binary log information for MySQL Cluster replication. See NDB Cluster Replication Schema and Tables.

  The `ndb_binlog_index` table uses the `MyISAM` storage engine. It is created only if the server is built with `NDB`.
- `slave_master_info`, `slave_relay_log_info`, `slave_worker_info`: Used to store replication information on slave servers. See Section 18.2.4, “Replication Relay and Status Logs”.

Optimizer System Tables

These system tables are for use by the optimizer:

- `innodb_index_stats`, `innodb_table_stats`: Used for `InnoDB` persistent optimizer statistics. See Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.
- `server_cost`, `engine_cost`: The optimizer cost model uses tables that contain cost estimate information about operations that occur during query execution. `server_cost` contains optimizer cost estimates for general server operations. `engine_cost` contains estimates for operations specific to particular storage engines. See Section 8.9.5, “The Optimizer Cost Model”.
- `column_stats`: Statistics about column values. In MySQL 8.0.2, this table was replaced by the `column_statistics` data dictionary table.

Miscellaneous System Tables

Other system tables do not fit the preceding categories:

- `servers`: Used by the `FEDERATED` storage engine. See Section 16.8.2.2, “Creating a FEDERATED Table Using CREATE SERVER”.
- `innodb_dynamic_metadata`: Used by the `InnoDB` storage engine to store fast-changing table metadata such as auto-increment counter values and index tree corruption flags. Replaces the data dictionary buffer table that resided in the `InnoDB` system tablespace. Introduced in MySQL 8.0.2.

5.4 MySQL Server Logs

MySQL Server has several logs that can help you find out what activity is taking place.

Log Type	Information Written to Log
Error log	Problems encountered starting, running, or stopping `mysqld`
General query log	Established client connections and statements received from clients
Binary log	Statements that change data (also used for replication)
Relay log	Data changes received from a replication master server
Slow query log	Queries that took more than `long_query_time` seconds to execute
DDL log (metadata log)	Metadata operations performed by DDL statements

By default, no logs are enabled, except the error log on Windows. (The DDL log is always created when required, and has no user-configurable options; see Section 5.4.6, “The DDL Log”.) The following log-specific sections provide information about the server options that enable logging.

By default, the server writes files for all enabled logs in the data directory. You can force the server to close and reopen the log files (or in some cases switch to a new log file) by flushing the logs. Log flushing occurs when you issue a `FLUSH LOGS` statement; execute `mysqladmin` with a `flush-logs` or `refresh` argument; or execute `mysqldump` with a `--flush-logs` or `--master-data` option. See Section 13.7.6.3, “FLUSH Syntax”, Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”, and Section 4.5.4, “`mysqldump` — A Database Backup Program”. In addition, the binary log is flushed when its size reaches the value of the `max_binlog_size` system variable.

You can control the general query and slow query logs during runtime. You can enable or disable logging, or change the log file name. You can tell the server to write general query and slow query entries to log tables, log files, or both. For details, see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”, Section 5.4.3, “The General Query Log”, and Section 5.4.5, “The Slow Query Log”.

The relay log is used only on slave replication servers, to hold data changes from the master server that must also be made on the slave. For discussion of relay log contents and configuration, see Section 18.2.4.1, “The Slave Relay Log”.

For information about log maintenance operations such as expiration of old log files, see Section 5.4.7, “Server Log Maintenance”.

For information about keeping logs secure, see Section 6.1.2.3, “Passwords and Logging”.

5.4.1 Selecting General Query and Slow Query Log Output Destinations

MySQL Server provides flexible control over the destination of output to the general query log and the slow query log, if those logs are enabled. Possible destinations for log entries are log files or the `general_log` and `slow_log` tables in the `mysql` database. Either or both destinations can be selected.

Log control at server startup. The `--log-output` option specifies the destination for log output. This option does not in itself enable the logs. Its syntax is `--log-output[=value,...]`:

- If `--log-output` is given with a value, the value should be a comma-separated list of one or more of the words `TABLE` (log to tables), `FILE` (log to files), or `NONE` (do not log to tables or files). `NONE`, if present, takes precedence over any other specifiers.
- If `--log-output` is omitted, the default logging destination is `FILE`.

The `general_log` system variable controls logging to the general query log for the selected log destinations. If specified at server startup, `general_log` takes an optional argument of 1 or 0 to enable or disable the log. To specify a file name other than the default for file logging, set the `general_log_file` variable. Similarly, the `slow_query_log` variable controls logging to the slow query log for the selected destinations and setting `slow_query_log_file` specifies a file name for file logging. If either log is enabled, the server opens the corresponding log file and writes startup messages to it. However, further logging of queries to the file does not occur unless the `FILE` log destination is selected.

Examples:

- To write general query log entries to the log table and the log file, use `--log-output=TABLE,FILE` to select both log destinations and `--general_log` to enable the general query log.
- To write general and slow query log entries only to the log tables, use `--log-output=TABLE` to select tables as the log destination and `--general_log` and `--slow_query_log` to enable both logs.

- To write slow query log entries only to the log file, use `--log-output=FILE` to select files as the log destination and `--slow_query_log` to enable the slow query log. (In this case, because the default log destination is `FILE`, you could omit the `--log-output` option.)

Log control at runtime. The system variables associated with log tables and files enable runtime control over logging:

- The global `log_output` system variable indicates the current logging destination. It can be modified at runtime to change the destination.
- The global `general_log` and `slow_query_log` variables indicate whether the general query log and slow query log are enabled (`ON`) or disabled (`OFF`). You can set these variables at runtime to control whether the logs are enabled.
- The global `general_log_file` and `slow_query_log_file` variables indicate the names of the general query log and slow query log files. You can set these variables at server startup or at runtime to change the names of the log files.
- To disable or enable general query logging for the current connection, set the session `sql_log_off` variable to `ON` or `OFF`.

The use of tables for log output offers the following benefits:

- Log entries have a standard format. To display the current structure of the log tables, use these statements:

```
SHOW CREATE TABLE mysql.general_log;
SHOW CREATE TABLE mysql.slow_log;
```

- Log contents are accessible through SQL statements. This enables the use of queries that select only those log entries that satisfy specific criteria. For example, to select log contents associated with a particular client (which can be useful for identifying problematic queries from that client), it is easier to do this using a log table than a log file.
- Logs are accessible remotely through any client that can connect to the server and issue queries (if the client has the appropriate log table privileges). It is not necessary to log in to the server host and directly access the file system.

The log table implementation has the following characteristics:

- In general, the primary purpose of log tables is to provide an interface for users to observe the runtime execution of the server, not to interfere with its runtime execution.
- `CREATE TABLE`, `ALTER TABLE`, and `DROP TABLE` are valid operations on a log table. For `ALTER TABLE` and `DROP TABLE`, the log table cannot be in use and must be disabled, as described later.
- By default, the log tables use the `CSV` storage engine that writes data in comma-separated values format. For users who have access to the `.CSV` files that contain log table data, the files are easy to import into other programs such as spreadsheets that can process CSV input.

  The log tables can be altered to use the `MyISAM` storage engine. You cannot use `ALTER TABLE` to alter a log table that is in use. The log must be disabled first. No engines other than `CSV` or `MyISAM` are legal for the log tables.
- To disable logging so that you can alter (or drop) a log table, you can use the following strategy. The example uses the general query log; the procedure for the slow query log is similar but uses the `slow_log` table and `slow_query_log` system variable.

```
SET @old_log_state = @@global.general_log;
SET GLOBAL general_log = 'OFF';
ALTER TABLE mysql.general_log ENGINE = MyISAM;
SET GLOBAL general_log = @old_log_state;
```

- `TRUNCATE TABLE` is a valid operation on a log table. It can be used to expire log entries.
- `RENAME TABLE` is a valid operation on a log table. You can atomically rename a log table (to perform log rotation, for example) using the following strategy:

```
USE mysql;
DROP TABLE IF EXISTS general_log2;
CREATE TABLE general_log2 LIKE general_log;
RENAME TABLE general_log TO general_log_backup, general_log2 TO general_log;
```

- `CHECK TABLE` is a valid operation on a log table.
- `LOCK TABLES` cannot be used on a log table.
- `INSERT`, `DELETE`, and `UPDATE` cannot be used on a log table. These operations are permitted only internally to the server itself.
- `FLUSH TABLES WITH READ LOCK` and the state of the `read_only` system variable have no effect on log tables. The server can always write to the log tables.
- Entries written to the log tables are not written to the binary log and thus are not replicated to slave servers.
- To flush the log tables or log files, use `FLUSH TABLES` or `FLUSH LOGS`, respectively.
- Partitioning of log tables is not permitted.
- A `mysqldump` dump includes statements to recreate those tables so that they are not missing after reloading the dump file. Log table contents are not dumped.

5.4.2 The Error Log

The error log contains information about `mysqld` startup and shutdown times, progress notes that occur during startup and shutdown, and critical errors that occur while the server is running. If `mysqld` notices a table that needs to be automatically checked or repaired, it writes a message to the error log.

On some operating systems, the error log contains a stack trace if `mysqld` exits abnormally. The trace can be used to determine where `mysqld` exited. See Section 28.5, “Debugging and Porting MySQL”.

Interpretation of error logging options differs somewhat for Windows and Unix systems, so be sure to use the information appropriate to your system.

`mysqld_safe` may write messages to the error log if it is used to start `mysqld`. For example, when `mysqld_safe` notices abnormal `mysqld` exits, it restarts `mysqld` and writes a `mysqld restarted` message to the error log.

In the following discussion, “console” means `stderr`, the standard error output. This is your terminal or console window unless the standard error output has been redirected.

- Error Logging on Windows
- Error Logging on Unix and Unix-Like Systems
- Using Syslog for the Error Log
- Error Log Verbosity
- Error Log Message Format
- Flushing and Renaming the Error Log File

Error Logging on Windows

On Windows, the `--log-error`, `--pid-file`, and `--console` options affect error logging:

- Without `--log-error`, `mysqld` writes error messages to the default log file.
- With `--log-error[=file_name]`, `mysqld` writes error messages to an error log file. If no file is named, `mysqld` writes to the default log file. If a file is named, `mysqld` writes to it, creating it in the data directory unless an absolute path name is given to specify a different directory.
- The default log file is *`host_name`*`.err` in the data directory, unless the `--pid-file` option is specified. In that case, the default name is the PID file base name with a suffix of `.err` in the data directory. This default is used if `--log-error` is not given, or is given without naming a log file.
- With `--console`, `mysqld` writes error messages to the console. `--log-error`, if given, is ignored and has no effect. If both options are present, their order does not matter: `--console` takes precedence and error messages go to the console.

In addition, on Windows, the server by default writes events and error messages to the Windows Event Log within the Application log.

- Entries marked as `Error`, `Warning`, and `Note` are written to the Event Log, but not informational messages such as information statements from individual storage engines.
- These log entries have a source of `MySQL`.
- Information written to the Windows Event Log can be controlled using the `log_syslog` system variable, as described later.

If the server writes error messages to the console, it sets the `log_error` system variable to `stderr`. Otherwise, `log_error` indicates the error log file name. On Windows, `--console` overrides use of an error log file and sends error messages to the console, so the server sets `log_error` to `stderr`. This occurs even if `--log-error` is also given.

Error Logging on Unix and Unix-Like Systems

On Unix and Unix-like systems, `mysqld` writes error log messages as follows:

- Without `--log-error`, `mysqld` writes error messages to the console.
- With `--log-error[=file_name]`, `mysqld` writes error messages to an error log file. If no file is named, `mysqld` writes to the default log file. If a file is named, `mysqld` writes to it, creating it in the data directory unless an absolute path name is given to specify a different directory.
- The default log file is *`host_name`*`.err` in the data directory. This default is used if `--log-error` is given without naming a log file.

Note

It is common for Yum or APT package installations to configure an error log location under `/var/log` with an entry like `log-error=/var/log/mysqld.log` in a server configuration file. Removing the file name from the entry causes the default log file to be used, which is written to the data directory.

If the server writes error messages to the console, it sets the `log_error` system variable to `stderr`. Otherwise, `log_error` indicates the error log file name.

If you specify `--log-error` in an option file in a `[mysqld]`, `[server]`, or `[mysqld_safe]` section, `mysqld_safe` finds and uses the option.

Using Syslog for the Error Log

It is possible to have `mysqld` write the error log to `syslog` on Unix and Unix-like systems, and to the Windows Event Log on Windows. To do so, use these system variables:

- `log_syslog`: Enable this variable to send the error log to `syslog`. In this case, the following system variables can also be used for finer control.

- `log_syslog_facility`: The default facility for `syslog` messages is `daemon`. Set this variable to specify a different facility.
- `log_syslog_include_pid`: Whether to include the server process ID in each line of `syslog` output.
- `log_syslog_tag`: This variable defines a tag to add to the server identifier (`mysqld`) in `syslog` messages. If defined, the tag is appended to the identifier with a leading hyphen.

On Unix and Unix-like systems, control of output to `syslog` is also available using `mysqld_safe`, which can capture server error output and pass it to `syslog`.

Note

Using `mysqld_safe` for `syslog` error logging is deprecated; you should use the server system variables instead.

`mysqld_safe` has three error-logging options, `--syslog`, `--skip-syslog`, and `--log-error`. The default with no logging options or with `--skip-syslog` is to use the default log file. To explicitly specify use of an error log file, specify `--log-error=file_name` to `mysqld_safe`, and `mysqld_safe` will arrange for `mysqld` to write messages to a log file. To use `syslog` instead, specify the `--syslog` option. For `syslog` output, a tag can be specified with `--syslog-tag=tag_val`; this is appended to the `mysqld` server identifier with a leading hyphen.

Error Log Verbosity

The `log_error_verbosity` system variable controls server verbosity for writing error, warning, and note messages to the error log. Permitted values are 1 (errors only), 2 (errors and warnings), 3 (errors, warnings, and notes), with a default of 3. If the value is greater than 2, the server logs aborted connections and access-denied errors for new connection attempts. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

Error Log Message Format

The `log_timestamps` system variable controls the timestamp time zone of messages written to the error log (as well as to general query log and slow query log files). Permitted values are `UTC` (the default) and `SYSTEM` (local system time zone).

The ID included in error log messages is that of the thread within `mysqld` responsible for writing the message. This indicates which part of the server produced the message, and is consistent with general query log and slow query log messages, which include the connection thread ID.

Flushing and Renaming the Error Log File

If you flush the logs using `FLUSH LOGS` or `mysqladmin flush-logs` and `mysqld` is writing the error log to a file (for example, if it was started with the `--log-error` option), the server closes and reopens the log file. To rename the file, do so manually before flushing. Flushing the logs then reopens a new file with the original file name. For example, to rename the file and create a new one, use the following commands (assuming a log file name of *`host_name`*`.err`):

```
shell> mv host_name.err host_name.err-old
shell> mysqladmin flush-logs
shell> mv host_name.err-old backup-directory
```

On Windows, use `rename` rather than `mv`.

If the location of the error file is not writable by the server, the log-flushing operation will fail to create a new log file. For example, on Linux, the server might write the error log as `/var/log/mysqld.log`, where `/var/log` is owned by `root` and not writable by `mysqld`. For information about handling this case, see Section 5.4.7, “Server Log Maintenance”.

If the server is not writing to a named file, no error log renaming occurs when the logs are flushed.

5.4.3 The General Query Log

The general query log is a general record of what `mysqld` is doing. The server writes information to this log when clients connect or disconnect, and it logs each SQL statement received from clients. The general query log can be very useful when you suspect an error in a client and want to know exactly what the client sent to `mysqld`.

Each line that shows when a client connects also includes `using` *`connection_type`* to indicate the protocol used to establish the connection. *`connection_type`* is one of `TCP/IP` (TCP/IP connection established without SSL), `SSL/TLS` (TCP/IP connection established with SSL), `Socket` (Unix socket file connection), `Named Pipe` (Windows named pipe connection), or `Shared Memory` (Windows shared memory connection).

`mysqld` writes statements to the query log in the order that it receives them, which might differ from the order in which they are executed. This logging order is in contrast with that of the binary log, for which statements are written after they are executed but before any locks are released. In addition, the query log may contain statements that only select data while such statements are never written to the binary log.

When using statement-based binary logging on a replication master server, statements received by its slaves are written to the query log of each slave. Statements are written to the query log of the master server if a client reads events with the `mysqlbinlog` utility and passes them to the server.

However, when using row-based binary logging, updates are sent as row changes rather than SQL statements, and thus these statements are never written to the query log when `binlog_format` is `ROW`. A given update also might not be written to the query log when this variable is set to `MIXED`, depending on the statement used. See Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”, for more information.

By default, the general query log is disabled. To specify the initial general query log state explicitly, use `--general_log[={0|1}]`. With no argument or an argument of 1, `--general_log` enables the log. With an argument of 0, this option disables the log. To specify a log file name, use `--general_log_file=`*`file_name`*. To specify the log destination, use `--log-output` (as described in Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”).

If you specify no name for the general query log file, the default name is *`host_name`*`.log`. The server creates the file in the data directory unless an absolute path name is given to specify a different directory.

To disable or enable the general query log or change the log file name at runtime, use the global `general_log` and `general_log_file` system variables. Set `general_log` to 0 (or `OFF`) to disable the log or to 1 (or `ON`) to enable it. Set `general_log_file` to specify the name of the log file. If a log file already is open, it is closed and the new file is opened.

When the general query log is enabled, the server writes output to any destinations specified by the `--log-output` option or `log_output` system variable. If you enable the log, the server opens the log file and writes startup messages to it. However, further logging of queries to the file does not occur unless the `FILE` log destination is selected. If the destination is `NONE`, the server writes no queries even if the general log is enabled. Setting the log file name has no effect on logging if the log destination value does not contain `FILE`.

Server restarts and log flushing do not cause a new general query log file to be generated (although flushing closes and reopens it). To rename the file and create a new one, use the following commands:

```
shell> mv host_name.log host_name-old.log
shell> mysqladmin flush-logs
shell> mv host_name-old.log backup-directory
```

On Windows, use `rename` rather than `mv`.

You can also rename the general query log file at runtime by disabling the log:

```
SET GLOBAL general_log = 'OFF';
```

With the log disabled, rename the log file externally; for example, from the command line. Then enable the log again:

```
SET GLOBAL general_log = 'ON';
```

This method works on any platform and does not require a server restart.

The session `sql_log_off` variable can be set to `ON` or `OFF` to disable or enable general query logging for the current connection.

Passwords in statements written to the general query log are rewritten by the server not to occur literally in plain text. Password rewriting can be suppressed for the general query log by starting the server with the `--log-raw` option. This option may be useful for diagnostic purposes, to see the exact text of statements as received by the server, but for security reasons is not recommended for production use. See also Section 6.1.2.3, “Passwords and Logging”.

An implication of password rewriting is that statements that cannot be parsed (due, for example, to syntax errors) are not written to the general query log because they cannot be known to be password free. Use cases that require logging of all statements including those with errors should use the `--log-raw` option, bearing in mind that this also bypasses password rewriting.

Password rewriting occurs only when plain text passwords are expected. For statements with syntax that expect a password hash value, no rewriting occurs. If a plain text password is supplied erroneously for such syntax, the password is logged as given, without rewriting. For example, the following statement is logged as shown because a password hash value is expected:

```
CREATE USER 'user1'@'localhost' IDENTIFIED BY PASSWORD 'not-so-secret';
```

The `log_timestamps` system variable controls the timestamp time zone of messages written to the general query log file (as well as to the slow query log file and the error log). It does not affect the time zone of general query log and slow query log messages written to log tables, but rows retrieved from those tables can be converted from the local system time zone to any desired time zone with `CONVERT_TZ()` or by setting the session `time_zone` system variable.

5.4.4 The Binary Log

The binary log contains “events” that describe database changes such as table creation operations or changes to table data. It also contains events for statements that potentially could have made changes (for example, a `DELETE` which matched no rows), unless row-based logging is used. The binary log also contains information about how long each statement took that updated data. The binary log has two important purposes:

- For replication, the binary log on a master replication server provides a record of the data changes to be sent to slave servers. The master server sends the events contained in its binary log to its slaves, which execute those events to make the same data changes that were made on the master. See Section 18.2, “Replication Implementation”.
- Certain data recovery operations require use of the binary log. After a backup has been restored, the events in the binary log that were recorded after the backup was made are re-executed. These events bring databases up to date from the point of the backup. See Section 7.5, “Point-in-Time (Incremental) Recovery Using the Binary Log”.

The binary log is not used for statements such as `SELECT` or `SHOW` that do not modify data. To log all statements (for example, to identify a problem query), use the general query log. See Section 5.4.3, “The General Query Log”.

Running a server with binary logging enabled makes performance slightly slower. However, the benefits of the binary log in enabling you to set up replication and for restore operations generally outweigh this minor performance decrement.

The binary log is crash-safe. Only complete events or transactions are logged or read back.

Passwords in statements written to the binary log are rewritten by the server not to occur literally in plain text. See also Section 6.1.2.3, “Passwords and Logging”.

The following discussion describes some of the server options and variables that affect the operation of binary logging. For a complete list, see Section 18.1.6.4, “Binary Logging Options and Variables”.

To enable the binary log, start the server with the `--log-bin[=base_name]` option. If no *`base_name`* value is given, the default name is the value of the `pid-file` option (which by default is the name of host machine) followed by `-bin`. If the base name is given, the server writes the file in the data directory unless the base name is given with a leading absolute path name to specify a different directory. It is recommended that you specify a base name explicitly rather than using the default of the host name; see Section B.5.7, “Known Issues in MySQL”, for the reason.

If you supply an extension in the log name (for example, `--log-bin=base_name.extension`), the extension is silently removed and ignored.

`mysqld` appends a numeric extension to the binary log base name to generate binary log file names. The number increases each time the server creates a new log file, thus creating an ordered series of files. The server creates a new file in the series each time it starts or flushes the logs. The server also creates a new binary log file automatically after the current log's size reaches `max_binlog_size`. A binary log file may become larger than `max_binlog_size` if you are using large transactions because a transaction is written to the file in one piece, never split between files.

To keep track of which binary log files have been used, `mysqld` also creates a binary log index file that contains the names of all used binary log files. By default, this has the same base name as the binary log file, with the extension `'.index'`. You can change the name of the binary log index file with the `--log-bin-index[=file_name]` option. You should not manually edit this file while `mysqld` is running; doing so would confuse `mysqld`.

The term “binary log file” generally denotes an individual numbered file containing database events. The term “binary log” collectively denotes the set of numbered binary log files plus the index file.

A client that has the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege can disable binary logging of its own statements by using a `SET sql_log_bin=0` statement. See Section 5.1.5, “Server System Variables”.

By default, the server logs the length of the event as well as the event itself and uses this to verify that the event was written correctly. You can also cause the server to write checksums for the events by setting the `binlog_checksum` system variable. When reading back from the binary log, the master uses the event length by default, but can be made to use checksums if available by enabling the `master_verify_checksum` system variable. The slave I/O thread also verifies events received from the master. You can cause the slave SQL thread to use checksums if available when reading from the relay log by enabling the `slave_sql_verify_checksum` system variable.

The format of the events recorded in the binary log is dependent on the binary logging format. Three format types are supported, row-based logging, statement-based logging and mixed-base logging. The binary logging format used depends on the MySQL version. For general descriptions of the logging formats, see Section 5.4.4.1, “Binary Logging Formats”. For detailed information about the format of the binary log, see MySQL Internals: The Binary Log.

The server evaluates the `--binlog-do-db` and `--binlog-ignore-db` options in the same way as it does the `--replicate-do-db` and `--replicate-ignore-db` options. For information about how this is done, see Section 18.2.5.1, “Evaluation of Database-Level Replication and Binary Logging Options”.

A replication slave server by default does not write to its own binary log any data modifications that are received from the replication master. To log these modifications, start the slave with the `--log-slave-updates` option in addition to the `--log-bin` option (see Section 18.1.6.3, "Replication Slave Options and Variables"). This is done when a slave is also to act as a master to other slaves in chained replication.

You can delete all binary log files with the `RESET MASTER` statement, or a subset of them with `PURGE BINARY LOGS`. See Section 13.7.6.6, "RESET Syntax", and Section 13.4.1.1, "PURGE BINARY LOGS Syntax".

If you are using replication, you should not delete old binary log files on the master until you are sure that no slave still needs to use them. For example, if your slaves never run more than three days behind, once a day you can execute `mysqladmin flush-logs` on the master and then remove any logs that are more than three days old. You can remove the files manually, but it is preferable to use `PURGE BINARY LOGS`, which also safely updates the binary log index file for you (and which can take a date argument). See Section 13.4.1.1, "PURGE BINARY LOGS Syntax".

You can display the contents of binary log files with the `mysqlbinlog` utility. This can be useful when you want to reprocess statements in the log for a recovery operation. For example, you can update a MySQL server from the binary log as follows:

```
shell> mysqlbinlog log_file | mysql -h server_name
```

`mysqlbinlog` also can be used to display replication slave relay log file contents because they are written using the same format as binary log files. For more information on the `mysqlbinlog` utility and how to use it, see Section 4.6.8, "`mysqlbinlog` — Utility for Processing Binary Log Files". For more information about the binary log and recovery operations, see Section 7.5, "Point-in-Time (Incremental) Recovery Using the Binary Log".

Binary logging is done immediately after a statement or transaction completes but before any locks are released or any commit is done. This ensures that the log is logged in commit order.

Updates to nontransactional tables are stored in the binary log immediately after execution.

Within an uncommitted transaction, all updates (`UPDATE`, `DELETE`, or `INSERT`) that change transactional tables such as `InnoDB` tables are cached until a `COMMIT` statement is received by the server. At that point, `mysqld` writes the entire transaction to the binary log before the `COMMIT` is executed.

Modifications to nontransactional tables cannot be rolled back. If a transaction that is rolled back includes modifications to nontransactional tables, the entire transaction is logged with a `ROLLBACK` statement at the end to ensure that the modifications to those tables are replicated.

When a thread that handles the transaction starts, it allocates a buffer of `binlog_cache_size` to buffer statements. If a statement is bigger than this, the thread opens a temporary file to store the transaction. The temporary file is deleted when the thread ends.

The `Binlog_cache_use` status variable shows the number of transactions that used this buffer (and possibly a temporary file) for storing statements. The `Binlog_cache_disk_use` status variable shows how many of those transactions actually had to use a temporary file. These two variables can be used for tuning `binlog_cache_size` to a large enough value that avoids the use of temporary files.

The `max_binlog_cache_size` system variable (default 4GB, which is also the maximum) can be used to restrict the total size used to cache a multiple-statement transaction. If a transaction is larger than this many bytes, it fails and rolls back. The minimum value is 4096.

If you are using the binary log and row based logging, concurrent inserts are converted to normal inserts for `CREATE ... SELECT` or `INSERT ... SELECT` statements. This is done to ensure that you can re-create an exact copy of your tables by applying the log during a backup operation. If you are using statement-based logging, the original statement is written to the log.

The binary log format has some known limitations that can affect recovery from backups. See Section 18.4.1, “Replication Features and Issues”.

Binary logging for stored programs is done as described in Section 23.7, “Binary Logging of Stored Programs”.

Note that the binary log format differs in MySQL 8.0 from previous versions of MySQL, due to enhancements in replication. See Section 18.4.2, “Replication Compatibility Between MySQL Versions”.

Writes to the binary log file and binary log index file are handled in the same way as writes to `MyISAM` tables. See Section B.5.3.4, “How MySQL Handles a Full Disk”.

By default, the binary log is synchronized to disk at each write (`sync_binlog=1`). If `sync_binlog` was not enabled, and the operating system or machine (not only the MySQL server) crashed, there is a chance that the last statements of the binary log could be lost. To prevent this, enable the `sync_binlog` system variable to synchronize the binary log to disk after every *`N`* commit groups. See Section 5.1.5, “Server System Variables”. The safest value for `sync_binlog` is 1 (the default), but this is also the slowest.

In earlier MySQL releases, there was a chance of inconsistency between the table content and binary log content if a crash occurred, even with `sync_binlog` set to 1. For example, if you are using `InnoDB` tables and the MySQL server processes a `COMMIT` statement, it writes many prepared transactions to the binary log in sequence, synchronizes the binary log, and then commits the transaction into `InnoDB`. If the server crashed between those two operations, the transaction would be rolled back by `InnoDB` at restart but still exist in the binary log. Such an issue was resolved in previous releases by enabling `InnoDB` support for two-phase commit in XA transactions. In 5.8.0 and higher, the `InnoDB` support for two-phase commit in XA transactions is always enabled.

`InnoDB` support for two-phase commit in XA transactions ensures that the binary log and `InnoDB` data files are synchronized. However, the MySQL server should also be configured to synchronize the binary log and the `InnoDB` logs to disk before committing the transaction. The `InnoDB` logs are synchronized by default, and `sync_binlog=1` ensures the binary log is synchronized. The effect of implicit `InnoDB` support for two-phase commit in XA transactions and `sync_binlog=1` is that at restart after a crash, after doing a rollback of transactions, the MySQL server scans the latest binary log file to collect transaction *`xid`* values and calculate the last valid position in the binary log file. The MySQL server then tells `InnoDB` to complete any prepared transactions that were successfully written to the to the binary log, and truncates the binary log to the last valid position. This ensures that the binary log reflects the exact data of `InnoDB` tables, and therefore the slave remains in synchrony with the master because it does not receive a statement which has been rolled back.

If the MySQL server discovers at crash recovery that the binary log is shorter than it should have been, it lacks at least one successfully committed `InnoDB` transaction. This should not happen if `sync_binlog=1` and the disk/file system do an actual sync when they are requested to (some do not), so the server prints an error message `The binary log file_name is shorter than its expected size`. In this case, this binary log is not correct and replication should be restarted from a fresh snapshot of the master's data.

The session values of the following system variables are written to the binary log and honored by the replication slave when parsing the binary log:

- `sql_mode` (except that the `NO_DIR_IN_CREATE` mode is not replicated; see Section 18.4.1.38, “Replication and Variables”)
- `foreign_key_checks`
- `unique_checks`
- `character_set_client`
- `collation_connection`

- `collation_database`
- `collation_server`
- `sql_auto_is_null`

5.4.4.1 Binary Logging Formats

The server uses several logging formats to record information in the binary log:

- Replication capabilities in MySQL originally were based on propagation of SQL statements from master to slave. This is called *statement-based logging*. You can cause this format to be used by starting the server with `--binlog-format=STATEMENT`.
- In *row-based logging* (the default), the master writes events to the binary log that indicate how individual table rows are affected. You can cause the server to use row-based logging by starting it with `--binlog-format=ROW`.
- A third option is also available: *mixed logging*. With mixed logging, statement-based logging is used by default, but the logging mode switches automatically to row-based in certain cases as described below. You can cause MySQL to use mixed logging explicitly by starting `mysqld` with the option `--binlog-format=MIXED`.

The logging format can also be set or limited by the storage engine being used. This helps to eliminate issues when replicating certain statements between a master and slave which are using different storage engines.

With statement-based replication, there may be issues with replicating nondeterministic statements. In deciding whether or not a given statement is safe for statement-based replication, MySQL determines whether it can guarantee that the statement can be replicated using statement-based logging. If MySQL cannot make this guarantee, it marks the statement as potentially unreliable and issues the warning, `Statement may not be safe to log in statement format`.

You can avoid these issues by using MySQL's row-based replication instead.

5.4.4.2 Setting The Binary Log Format

You can select the binary logging format explicitly by starting the MySQL server with `--binlog-format=type`. The supported values for *`type`* are:

- `STATEMENT` causes logging to be statement based.
- `ROW` causes logging to be row based. This is the default.
- `MIXED` causes logging to use mixed format.

The logging format also can be switched at runtime. To specify the format globally for all clients, set the global value of the `binlog_format` system variable:

```
mysql> SET GLOBAL binlog_format = 'STATEMENT';
mysql> SET GLOBAL binlog_format = 'ROW';
mysql> SET GLOBAL binlog_format = 'MIXED';
```

An individual client can control the logging format for its own statements by setting the session value of `binlog_format`:

```
mysql> SET SESSION binlog_format = 'STATEMENT';
mysql> SET SESSION binlog_format = 'ROW';
mysql> SET SESSION binlog_format = 'MIXED';
```

Note

Each MySQL Server can set its own and only its own binary logging format (true whether `binlog_format` is set with global or session scope). This means

> that changing the logging format on a replication master does not cause a slave to change its logging format to match. (When using `STATEMENT` mode, the `binlog_format` system variable is not replicated; when using `MIXED` or `ROW` logging mode, it is replicated but is ignored by the slave.) Changing the binary logging format on the master while replication is ongoing, or without also changing it on the slave can cause replication to fail with errors such as `Error executing row event: 'Cannot execute statement: impossible to write to binary log since statement is in row format and BINLOG_FORMAT = STATEMENT.'`

To change the global or session `binlog_format` value, you must have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege.

There are several reasons why a client might want to set binary logging on a per-session basis:

- A session that makes many small changes to the database might want to use row-based logging.
- A session that performs updates that match many rows in the `WHERE` clause might want to use statement-based logging because it will be more efficient to log a few statements than many rows.
- Some statements require a lot of execution time on the master, but result in just a few rows being modified. It might therefore be beneficial to replicate them using row-based logging.

There are exceptions when you cannot switch the replication format at runtime:

- From within a stored function or a trigger
- If the session is currently in row-based replication mode and has open temporary tables

Trying to switch the format in any of these cases results in an error.

If you are using `InnoDB` tables and the transaction isolation level is `READ COMMITTED` or `READ UNCOMMITTED`, only row-based logging can be used. It is *possible* to change the logging format to `STATEMENT`, but doing so at runtime leads very rapidly to errors because `InnoDB` can no longer perform inserts.

Switching the replication format at runtime is not recommended when any temporary tables exist, because temporary tables are logged only when using statement-based replication, whereas with row-based replication they are not logged. With mixed replication, temporary tables are usually logged; exceptions happen with user-defined functions (UDFs) and with the `UUID()` function.

With the binary log format set to `ROW`, many changes are written to the binary log using the row-based format. Some changes, however, still use the statement-based format. Examples include all DDL (data definition language) statements such as `CREATE TABLE`, `ALTER TABLE`, or `DROP TABLE`.

The `--binlog-row-event-max-size` option is available for servers that are capable of row-based replication. Rows are stored into the binary log in chunks having a size in bytes not exceeding the value of this option. The value must be a multiple of 256. The default value is 8192.

> **Warning**
>
> When using *statement-based logging* for replication, it is possible for the data on the master and slave to become different if a statement is designed in such a way that the data modification is *nondeterministic*; that is, it is left to the will of the query optimizer. In general, this is not a good practice even outside of replication. For a detailed explanation of this issue, see Section B.5.7, “Known Issues in MySQL”.

For information about logs kept by replication slaves, see Section 18.2.4, “Replication Relay and Status Logs”.

5.4.4.3 Mixed Binary Logging Format

When running in `MIXED` logging format, the server automatically switches from statement-based to row-based logging under the following conditions:

- When a function contains `UUID()`.
- When one or more tables with `AUTO_INCREMENT` columns are updated and a trigger or stored function is invoked. Like all other unsafe statements, this generates a warning if `binlog_format = STATEMENT`.

  For more information, see Section 18.4.1.1, “Replication and AUTO_INCREMENT”.
- When the body of a view requires row-based replication, the statement creating the view also uses it. For example, this occurs when the statement creating a view uses the `UUID()` function.
- When a call to a UDF is involved.
- If a statement is logged by row and the session that executed the statement has any temporary tables, logging by row is used for all subsequent statements (except for those accessing temporary tables) until all temporary tables in use by that session are dropped.

  This is true whether or not any temporary tables are actually logged.

  Temporary tables cannot be logged using row-based format; thus, once row-based logging is used, all subsequent statements using that table are unsafe. The server approximates this condition by treating all statements executed during the session as unsafe until the session no longer holds any temporary tables.
- When `FOUND_ROWS()` or `ROW_COUNT()` is used. (Bug #12092, Bug #30244)
- When `USER()`, `CURRENT_USER()`, or `CURRENT_USER` is used. (Bug #28086)
- When a statement refers to one or more system variables. (Bug #31168)

  Exception. The following system variables, when used with session scope (only), do not cause the logging format to switch:

  - `auto_increment_increment`
  - `auto_increment_offset`
  - `character_set_client`
  - `character_set_connection`
  - `character_set_database`
  - `character_set_server`
  - `collation_connection`
  - `collation_database`
  - `collation_server`
  - `foreign_key_checks`
  - `identity`
  - `last_insert_id`
  - `lc_time_names`

- `pseudo_thread_id`
- `sql_auto_is_null`
- `time_zone`
- `timestamp`
- `unique_checks`

For information about determining system variable scope, see Section 5.1.6, “Using System Variables”.

For information about how replication treats `sql_mode`, see Section 18.4.1.38, “Replication and Variables”.

- When one of the tables involved is a log table in the `mysql` database.
- When the `LOAD_FILE()` function is used. (Bug #39701)

Note

A warning is generated if you try to execute a statement using statement-based logging that should be written using row-based logging. The warning is shown both in the client (in the output of `SHOW WARNINGS`) and through the `mysqld` error log. A warning is added to the `SHOW WARNINGS` table each time such a statement is executed. However, only the first statement that generated the warning for each client session is written to the error log to prevent flooding the log.

In addition to the decisions above, individual engines can also determine the logging format used when information in a table is updated. The logging capabilities of an individual engine can be defined as follows:

- If an engine supports row-based logging, the engine is said to be *row-logging capable*.
- If an engine supports statement-based logging, the engine is said to be *statement-logging capable*.

A given storage engine can support either or both logging formats. The following table lists the formats supported by each engine.

Storage Engine	Row Logging Supported	Statement Logging Supported
`ARCHIVE`	Yes	Yes
`BLACKHOLE`	Yes	Yes
`CSV`	Yes	Yes
`EXAMPLE`	Yes	No
`FEDERATED`	Yes	Yes
`HEAP`	Yes	Yes
`InnoDB`	Yes	Yes when the transaction isolation level is `REPEATABLE READ` or `SERIALIZABLE`; No otherwise.
`MyISAM`	Yes	Yes
`MERGE`	Yes	Yes

Storage Engine	Row Logging Supported	Statement Logging Supported
NDB	Yes	No

Whether a statement is to be logged and the logging mode to be used is determined according to the type of statement (safe, unsafe, or binary injected), the binary logging format (STATEMENT, ROW, or MIXED), and the logging capabilities of the storage engine (statement capable, row capable, both, or neither). (Binary injection refers to logging a change that must be logged using ROW format.)

Statements may be logged with or without a warning; failed statements are not logged, but generate errors in the log. This is shown in the following decision table, where **SLC** stands for “statement-logging capable” and **RLC** stands for “row-logging capable”.

Condition				Action	
Type	**binlog_format**	**SLC**	**RLC**	**Error / Warning**	**Logged as**
*	*	No	No	Error: Cannot execute statement: Binary logging is impossible since at least one engine is involved that is both row-incapable and statement-incapable.	-
Safe	STATEMENT	Yes	No	-	STATEMENT
Safe	MIXED	Yes	No	-	STATEMENT
Safe	ROW	Yes	No	Error: Cannot execute statement: Binary logging is impossible since BINLOG_FORMAT = ROW and at least one table uses a storage engine that is not capable of row-based logging.	-
Unsafe	STATEMENT	Yes	No	Warning: Unsafe statement binlogged in statement format, since BINLOG_FORMAT = STATEMENT	STATEMENT
Unsafe	MIXED	Yes	No	Error: Cannot execute statement: Binary logging of an unsafe statement is impossible when the storage engine is limited to	-

Condition				Action	
Type	**binlog_format**	**SLC**	**RLC**	**Error / Warning**	**Logged as**
				statement-based logging, even if BINLOG_FORMAT = MIXED.	
Unsafe	ROW	Yes	No	Error: Cannot execute statement: Binary logging is impossible since BINLOG_FORMAT = ROW and at least one table uses a storage engine that is not capable of row-based logging.	-
Row Injection	STATEMENT	Yes	No	Error: Cannot execute row injection: Binary logging is not possible since at least one table uses a storage engine that is not capable of row-based logging.	-
Row Injection	MIXED	Yes	No	Error: Cannot execute row injection: Binary logging is not possible since at least one table uses a storage engine that is not capable of row-based logging.	-
Row Injection	ROW	Yes	No	Error: Cannot execute row injection: Binary logging is not possible since at least one table uses a storage engine that is not capable of row-based logging.	-
Safe	STATEMENT	No	Yes	Error: Cannot execute statement: Binary logging is impossible since BINLOG_FORMAT = STATEMENT and at least one table	-

Condition				Action	
Type	**binlog_format**	**SLC**	**RLC**	**Error / Warning**	**Logged as**
				uses a storage engine that is not capable of statement-based logging.	
Safe	MIXED	No	Yes	-	ROW
Safe	ROW	No	Yes	-	ROW
Unsafe	STATEMENT	No	Yes	Error: Cannot execute statement: Binary logging is impossible since BINLOG_FORMAT = STATEMENT and at least one table uses a storage engine that is not capable of statement-based logging.	-
Unsafe	MIXED	No	Yes	-	ROW
Unsafe	ROW	No	Yes	-	ROW
Row Injection	STATEMENT	No	Yes	Error: Cannot execute row injection: Binary logging is not possible since BINLOG_FORMAT = STATEMENT.	-
Row Injection	MIXED	No	Yes	-	ROW
Row Injection	ROW	No	Yes	-	ROW
Safe	STATEMENT	Yes	Yes	-	STATEMENT
Safe	MIXED	Yes	Yes	-	STATEMENT
Safe	ROW	Yes	Yes	-	ROW
Unsafe	STATEMENT	Yes	Yes	Warning: Unsafe statement binlogged in statement format since BINLOG_FORMAT = STATEMENT.	STATEMENT
Unsafe	MIXED	Yes	Yes	-	ROW
Unsafe	ROW	Yes	Yes	-	ROW
Row Injection	STATEMENT	Yes	Yes	Error: Cannot execute row injection:	-

Condition				Action	
Type	**binlog_format**	**SLC**	**RLC**	**Error / Warning**	**Logged as**
				Binary logging is not possible because BINLOG_FORMAT = STATEMENT.	
Row Injection	MIXED	Yes	Yes	-	ROW
Row Injection	ROW	Yes	Yes	-	ROW

When a warning is produced by the determination, a standard MySQL warning is produced (and is available using SHOW WARNINGS). The information is also written to the mysqld error log. Only one error for each error instance per client connection is logged to prevent flooding the log. The log message includes the SQL statement that was attempted.

If a slave server was started with log_error_verbosity set to display warnings, the slave prints messages to the error log to provide information about its status, such as the binary log and relay log coordinates where it starts its job, when it is switching to another relay log, when it reconnects after a disconnect, statements that are unsafe for statement-based logging, and so forth.

5.4.4.4 Logging Format for Changes to mysql Database Tables

The contents of the grant tables in the mysql database can be modified directly (for example, with INSERT or DELETE) or indirectly (for example, with GRANT or CREATE USER). Statements that affect mysql database tables are written to the binary log using the following rules:

- Data manipulation statements that change data in mysql database tables directly are logged according to the setting of the binlog_format system variable. This pertains to statements such as INSERT, UPDATE, DELETE, REPLACE, DO, LOAD DATA INFILE, SELECT, and TRUNCATE TABLE.
- Statements that change the mysql database indirectly are logged as statements regardless of the value of binlog_format. This pertains to statements such as GRANT, REVOKE, SET PASSWORD, RENAME USER, CREATE (all forms except CREATE TABLE ... SELECT), ALTER (all forms), and DROP (all forms).

CREATE TABLE ... SELECT is a combination of data definition and data manipulation. The CREATE TABLE part is logged using statement format and the SELECT part is logged according to the value of binlog_format.

5.4.5 The Slow Query Log

The slow query log consists of SQL statements that took more than long_query_time seconds to execute and required at least min_examined_row_limit rows to be examined. The minimum and default values of long_query_time are 0 and 10, respectively. The value can be specified to a resolution of microseconds. For logging to a file, times are written including the microseconds part. For logging to tables, only integer times are written; the microseconds part is ignored.

By default, administrative statements are not logged, nor are queries that do not use indexes for lookups. This behavior can be changed using log_slow_admin_statements and log_queries_not_using_indexes, as described later.

The time to acquire the initial locks is not counted as execution time. mysqld writes a statement to the slow query log after it has been executed and after all locks have been released, so log order might differ from execution order.

By default, the slow query log is disabled. To specify the initial slow query log state explicitly, use --slow_query_log[={0|1}]. With no argument or an argument of 1, --slow_query_log

enables the log. With an argument of 0, this option disables the log. To specify a log file name, use `--slow_query_log_file=file_name`. To specify the log destination, use `--log-output` (as described in Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”).

If you specify no name for the slow query log file, the default name is *`host_name`*`-slow.log`. The server creates the file in the data directory unless an absolute path name is given to specify a different directory.

To disable or enable the slow query log or change the log file name at runtime, use the global `slow_query_log` and `slow_query_log_file` system variables. Set `slow_query_log` to 0 (or `OFF`) to disable the log or to 1 (or `ON`) to enable it. Set `slow_query_log_file` to specify the name of the log file. If a log file already is open, it is closed and the new file is opened.

When the slow query log is enabled, the server writes output to any destinations specified by the `--log-output` option or `log_output` system variable. If you enable the log, the server opens the log file and writes startup messages to it. However, further logging of queries to the file does not occur unless the `FILE` log destination is selected. If the destination is `NONE`, the server writes no queries even if the slow query log is enabled. Setting the log file name has no effect on logging if the log destination value does not contain `FILE`.

The server writes less information to the slow query log if you use the `--log-short-format` option.

To include slow administrative statements in the statements written to the slow query log, use the `log_slow_admin_statements` system variable. Administrative statements include `ALTER TABLE`, `ANALYZE TABLE`, `CHECK TABLE`, `CREATE INDEX`, `DROP INDEX`, `OPTIMIZE TABLE`, and `REPAIR TABLE`.

To include queries that do not use indexes for row lookups in the statements written to the slow query log, enable the `log_queries_not_using_indexes` system variable. When such queries are logged, the slow query log may grow quickly. It is possible to put a rate limit on these queries by setting the `log_throttle_queries_not_using_indexes` system variable. By default, this variable is 0, which means there is no limit. Positive values impose a per-minute limit on logging of queries that do not use indexes. The first such query opens a 60-second window within which the server logs queries up to the given limit, then suppresses additional queries. If there are suppressed queries when the window ends, the server logs a summary that indicates how many there were and the aggregate time spent in them. The next 60-second window begins when the server logs the next query that does not use indexes.

The server uses the controlling parameters in the following order to determine whether to write a query to the slow query log:

1. The query must either not be an administrative statement, or `log_slow_admin_statements` must be enabled.

2. The query must have taken at least `long_query_time` seconds, or `log_queries_not_using_indexes` must be enabled and the query used no indexes for row lookups.

3. The query must have examined at least `min_examined_row_limit` rows.

4. The query must not be suppressed according to the `log_throttle_queries_not_using_indexes` setting.

The `log_timestamps` system variable controls the timestamp time zone of messages written to the slow query log file (as well as to the general query log file and the error log). It does not affect the time zone of general query log and slow query log messages written to log tables, but rows retrieved from those tables can be converted from the local system time zone to any desired time zone with `CONVERT_TZ()` or by setting the session `time_zone` system variable.

All log lines contain a timestamp.

The server does not write queries handled by the query cache to the slow query log, nor queries that would not benefit from the presence of an index because the table has zero rows or one row.

By default, a replication slave does not write replicated queries to the slow query log. To change this, use the `log_slow_slave_statements` system variable.

Passwords in statements written to the slow query log are rewritten by the server not to occur literally in plain text. See also Section 6.1.2.3, “Passwords and Logging”.

The slow query log can be used to find queries that take a long time to execute and are therefore candidates for optimization. However, examining a long slow query log can become a difficult task. To make this easier, you can process a slow query log file using the `mysqldumpslow` command to summarize the queries that appear in the log. See Section 4.6.9, “`mysqldumpslow` — Summarize Slow Query Log Files”.

5.4.6 The DDL Log

The DDL log, or metadata log, records metadata operations generated by data definition statements such as `DROP TABLE` and `ALTER TABLE`. MySQL uses this log to recover from crashes occurring in the middle of a metadata operation. When executing the statement `DROP TABLE t1, t2`, we need to ensure that both `t1` and `t2` are dropped, and that each table drop is complete. Another example of this type of SQL statement is `ALTER TABLE t3 DROP PARTITION p2`, where we must make certain that the partition is completely dropped and that its definition is removed from the list of partitions for table `t3`.

A record of metadata operations such as those just described are written to the file `ddl_log.log`, in the MySQL data directory. This is a binary file; it is not intended to be human-readable, and you should not attempt to modify its contents in any way.

`ddl_log.log` is not created until it is actually needed for recording metadata statements, and is removed following a successful start of `mysqld`. Thus, it is possible for this file not to be present on a MySQL server that is functioning in a completely normal manner.

Currently, `ddl_log.log` can hold up to 1048573 entries, equivalent 4 GB in size. Once this limit is exceeded, you must rename or remove the file before it is possible to execute any additional DDL statements. This is a known issue which we are working to resolve (Bug #83708).

There are no user-configurable server options or variables associated with this file.

5.4.7 Server Log Maintenance

As described in Section 5.4, “MySQL Server Logs”, MySQL Server can create several different log files to help you see what activity is taking place. However, you must clean up these files regularly to ensure that the logs do not take up too much disk space.

When using MySQL with logging enabled, you may want to back up and remove old log files from time to time and tell MySQL to start logging to new files. See Section 7.2, “Database Backup Methods”.

On a Linux (Red Hat) installation, you can use the `mysql-log-rotate` script for this. If you installed MySQL from an RPM distribution, this script should have been installed automatically. Be careful with this script if you are using the binary log for replication. You should not remove binary logs until you are certain that their contents have been processed by all slaves.

On other systems, you must install a short script yourself that you start from `cron` (or its equivalent) for handling log files.

For the binary log, you can set the `expire_logs_days` system variable to expire binary log files automatically after a given number of days (see Section 5.1.5, “Server System Variables”). If you are using replication, you should set the variable no lower than the maximum number of days your

slaves might lag behind the master. To remove binary logs on demand, use the `PURGE BINARY LOGS` statement (see Section 13.4.1.1, “PURGE BINARY LOGS Syntax”).

You can force MySQL to start using new log files by flushing the logs. Log flushing occurs when you issue a `FLUSH LOGS` statement or execute a `mysqladmin flush-logs`, `mysqladmin refresh`, `mysqldump --flush-logs`, or `mysqldump --master-data` command. See Section 13.7.6.3, “FLUSH Syntax”, Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”, and Section 4.5.4, “`mysqldump` — A Database Backup Program”. In addition, the binary log is flushed when its size reaches the value of the `max_binlog_size` system variable.

`FLUSH LOGS` supports optional modifiers to enable selective flushing of individual logs (for example, `FLUSH BINARY LOGS`).

A log-flushing operation does the following:

- If general query logging or slow query logging to a log file is enabled, the server closes and reopens the general query log file or slow query log file.
- If binary logging is enabled, the server closes the current binary log file and opens a new log file with the next sequence number.
- If the server was started with the `--log-error` option to cause the error log to be written to a file, the server closes and reopens the log file.

The server creates a new binary log file when you flush the logs. However, it just closes and reopens the general and slow query log files. To cause new files to be created on Unix, rename the current log files before flushing them. At flush time, the server opens new log files with the original names. For example, if the general and slow query log files are named `mysql.log` and `mysql-slow.log`, you can use a series of commands like this:

```
shell> cd mysql-data-directory
shell> mv mysql.log mysql.old
shell> mv mysql-slow.log mysql-slow.old
shell> mysqladmin flush-logs
```

On Windows, use `rename` rather than `mv`.

At this point, you can make a backup of `mysql.old` and `mysql-slow.old` and then remove them from disk.

A similar strategy can be used to back up the error log file, if there is one.

You can rename the general query log or slow query log at runtime by disabling the log:

```
SET GLOBAL general_log = 'OFF';
SET GLOBAL slow_query_log = 'OFF';
```

With the logs disabled, rename the log files externally; for example, from the command line. Then enable the logs again:

```
SET GLOBAL general_log = 'ON';
SET GLOBAL slow_query_log = 'ON';
```

This method works on any platform and does not require a server restart.

Note

For the server to recreate a given log file after you have renamed the file externally, the file location must be writable by the server. This may not always be the case. For example, on Linux, the server might write the error log as `/`

> `var/log/mysqld.log`, where `/var/log` is owned by `root` and not writable by `mysqld`. In this case, the log-flushing operation will fail to create a new log file.
>
> To handle this situation, you must manually create the new log file with the proper ownershiop after renaming the original log file. For example, execute these commands as `root`:
>
> ```
> shell> mv /var/log/mysqld.log /var/log/mysqld.log.old
> shell> install -omysql -gmysql -m0644 /dev/null /var/log/mysqld.log
> ```

5.5 MySQL Server Components

MySQL Server includes a component-based infrastructure for improving server extensibility:

- A component provides services that are available to the server and other components. (With respect to service use, the server is a component, equal to other components.) Components interact with each other only through the services they provide.
- Components are installed and uninstalled using the `INSTALL COMPONENT` and `UNINSTALL COMPONENT` SQL statements.
- A loader service handles component installation and uninstallation, and also lists installed components in the `mysql.component` system table, which serves as a registry.

The SQL statements for component manipulation affect server operation and the `mysql.component` system table as follows:

- `INSTALL COMPONENT` installs components into the server. The components become active immediately. The loader service also registers installed components in the `mysql.component` table. For subsequent server restarts, any components listed in `mysql.component` are installed by the loader service during the startup sequence. This occurs even if the server is started with the `--skip-grant-tables` option.
- `UNINSTALL COMPONENT` deactivates components and uninstalls them from the server. The loader service also unregisters the components from the `mysql.component` table so that they are no longer installed during the startup sequence for subsequent server restarts.

For information about the internal implementation of components, see http://dev.mysql.com/doc/dev/mysql-server/latest/PAGE_EXTENDING.html. For example, if you intend to write your own components, this information is important for understanding how components work.

5.6 MySQL Server Plugins

MySQL supports a plugin API that enables creation of server components. Plugins can be loaded at server startup, or loaded and unloaded at runtime without restarting the server. The components supported by this interface include, but are not limited to, storage engines, `INFORMATION_SCHEMA` tables, full-text parser plugins, and server extensions.

5.6.1 Server Plugins Available

MySQL distributions include several plugins that implement server extensions:

- Plugins for authenticating attempts by clients to connect to MySQL Server. Plugins are available for several authentication protocols. See Section 6.3.10, “Pluggable Authentication”.
- A connection-control plugin that enables administrators to introduce an increasing delay after a certain number of consecutive failed client connection attempts. See Section 6.5.2, “The Connection-Control Plugins”.

- A password-validation plugin implements password strength policies and assesses the strength of potential passwords. See Section 6.5.3, “The Password Validation Plugin”.
- Semisynchronous replication plugins implement an interface to replication capabilities that permit the master to proceed as long as at least one slave has responded to each transaction. See Section 18.3.10, “Semisynchronous Replication”.
- A query rewrite plugin examines statements received by MySQL Server and possibly rewrites them before the server executes them. See Section 5.6.4, “The Rewriter Query Rewrite Plugin”
- Version Tokens enables creation of and synchronization around server tokens that applications can use to prevent accessing incorrect or out-of-date data. Version Tokens is based on a plugin library that implements a `version_tokens` plugin and a set of user-defined functions. See Section 5.6.5, “Version Tokens”.
- A keyring plugin provides secure storage for sensitive information. See Section 6.5.4, “The MySQL Keyring”.
- Test framework plugins test server services. For information about these plugins, see Plugins for Testing Plugin Services.

The following sections describe how to install and uninstall plugins, and how to determine at runtime which plugins are installed and obtain information about them. For information about writing plugins, see Section 28.2, “The MySQL Plugin API”.

5.6.2 Installing and Uninstalling Plugins

Server plugins must be loaded into the server before they can be used. MySQL supports plugin loading at server startup and runtime. It is also possible to control the activation state of loaded plugins at startup, and to unload them at runtime.

- Installing plugins
- Controlling plugin activation state
- Uninstalling plugins

While a plugin is loaded, information about it is available at runtime from the `INFORMATION_SCHEMA.PLUGINS` table and the `SHOW PLUGINS` statement. See Section 5.6.3, “Obtaining Server Plugin Information”.

Installing Plugins

Before a server plugin can be used, it must be installed using one of the following methods. In the descriptions, *`plugin_name`* stands for a plugin name such as `innodb`, `csv`, or `validate_password`.

Built-in plugins:

A built-in plugin is known by the server automatically. Normally, the server enables the plugin at startup. Some built-in plugins permit this to be changed with the `--plugin_name[=activation_state]` option.

Plugins registered in the `mysql.plugin` system table:

The `mysql.plugin` table serves as a registry of plugins (other than built-in plugins, which need not be registered). At startup, the server loads each plugin listed in the table. Normally, for a plugin loaded from the `mysql.plugin` table, the server also enables the plugin. This can be changed with the `--plugin_name[=activation_state]` option.

If the server is started with the `--skip-grant-tables` option, it does not consult the `mysql.plugin` table and does not load the plugins listed there.

Plugins named with command-line options:

A plugin located in a plugin library file can be loaded at server startup with the `--plugin-load`, `--plugin-load-add`, or `--early-plugin-load` option. Normally, for a plugin loaded at startup, the server also enables the plugin. This can be changed with the `--plugin_name[=activation_state]` option.

The `--plugin-load` and `--plugin-load-add` options load plugins after built-in plugins and storage engines have initialized during the server startup sequence. The `--early-plugin-load` option is used to load plugins that must be available prior to initialization of built-in plugins and storage engines.

The value of each plugin-loading option is a semicolon-separated list of *`name=plugin_library`* and *`plugin_library`* values. Each *`name`* is the name of a plugin to load, and *`plugin_library`* is the name of the library file that contains the plugin code. If a plugin library is named without any preceding plugin name, the server loads all plugins in the library. The server looks for plugin library files in the directory named by the `plugin_dir` system variable.

Plugin-loading options do not register any plugin in the `mysql.plugin` table. For subsequent restarts, the server loads the plugin again only if `--plugin-load`, `--plugin-load-add`, or `--early-plugin-load` is given again. That is, the option produces a one-time plugin-installation operation that persists for a single server invocation.

`--plugin-load`, `--plugin-load-add`, and `--early-plugin-load` enable plugins to be loaded even when `--skip-grant-tables` is given (which causes the server to ignore the `mysql.plugin` table). `--plugin-load`, `--plugin-load-add`, and `--early-plugin-load` also enable plugins to be loaded at startup that cannot be loaded at runtime.

The `--plugin-load-add` option complements the `--plugin-load` option:

- Each instance of `--plugin-load` resets the set of plugins to load at startup, whereas `--plugin-load-add` adds a plugin or plugins to the set of plugins to be loaded without resetting the current set. Consequently, if multiple instances of `--plugin-load` are specified, only the last one takes effect. With multiple instances of `--plugin-load-add`, all of them take effect.
- The argument format is the same as for `--plugin-load`, but multiple instances of `--plugin-load-add` can be used to avoid specifying a large set of plugins as a single long unwieldy `--plugin-load` argument.
- `--plugin-load-add` can be given in the absence of `--plugin-load`, but any instance of `--plugin-load-add` that appears before `--plugin-load` has no effect because `--plugin-load` resets the set of plugins to load.

For example, these options:

```
--plugin-load=x --plugin-load-add=y
```

are equivalent to this option:

```
--plugin-load="x;y"
```

But these options:

```
--plugin-load-add=y --plugin-load=x
```

are equivalent to this option:

```
```

```
--plugin-load=x
```

Plugins installed with the `INSTALL PLUGIN` statement:

A plugin located in a plugin library file can be loaded at runtime with the `INSTALL PLUGIN` statement. The statement also registers the plugin in the `mysql.plugin` table to cause the server to load it on subsequent restarts. For this reason, `INSTALL PLUGIN` requires the `INSERT` privilege for the `mysql.plugin` table.

The plugin library file base name depends on your platform. Common suffixes are `.so` for Unix and Unix-like systems, `.dll` for Windows.

Example: The `--plugin-load` option installs a plugin at server startup. To install a plugin named `myplugin` from a plugin library file named `somepluglib.so`, use these lines in a `my.cnf` file:

```
[mysqld]
plugin-load=myplugin=somepluglib.so
```

In this case, the plugin is not registered in `mysql.plugin`. Restarting the server without the `--plugin-load` option causes the plugin not to be loaded at startup.

Alternatively, the `INSTALL PLUGIN` statement causes the server to load the plugin code from the library file at runtime:

```
INSTALL PLUGIN myplugin SONAME 'somepluglib.so';
```

`INSTALL PLUGIN` also causes “permanent” plugin registration: The plugin is listed in the `mysql.plugin` table to ensure that the server loads it on subsequent restarts.

Many plugins can be loaded either at server startup or at runtime. However, if a plugin is designed such that it must be loaded and initialized during server startup, attempts to load it at runtime using `INSTALL PLUGIN` produce an error:

```
mysql> INSTALL PLUGIN myplugin SONAME 'somepluglib.so';
ERROR 1721 (HY000): Plugin 'myplugin' is marked as not dynamically
installable. You have to stop the server to install it.
```

In this case, you must use `--plugin-load`, `--plugin-load-add`, or `--early-plugin-load`.

If a plugin is named both using a `--plugin-load`, `--plugin-load-add`, or `--early-plugin-load` option and (as a result of an earlier `INSTALL PLUGIN` statement) in the `mysql.plugin` table, the server starts but writes these messages to the error log:

```
[ERROR] Function 'plugin_name' already exists
[Warning] Couldn't load plugin named 'plugin_name'
with soname 'plugin_object_file'.
```

Controlling Plugin Activation State

If the server knows about a plugin when it starts (for example, because the plugin is named using a `--plugin-load` option or is registered in the `mysql.plugin` table), the server loads and enables the plugin by default. It is possible to control activation state for such a plugin using a `--plugin_name[=activation_state]` startup option, where *`plugin_name`* is the name of the plugin to affect, such as `innodb`, `csv`, or `validate_password`. As with other options, dashes and underscores are interchangeable in option names. Also, activation state values are not case sensitive. For example, `--my_plugin=ON` and `--my-plugin=on` are equivalent.

- `--plugin_name=OFF`

Tells the server to disable the plugin. This may not be possible for certain built-in plugins, such as `mysql_native_password`.

- `--plugin_name[=ON]`

  Tells the server to enable the plugin. (Specifying the option as `--plugin_name` without a value has the same effect.) If the plugin fails to initialize, the server runs with the plugin disabled.

- `--plugin_name=FORCE`

  Tells the server to enable the plugin, but if plugin initialization fails, the server does not start. In other words, this option forces the server to run with the plugin enabled or not at all.

- `--plugin_name=FORCE_PLUS_PERMANENT`

  Like `FORCE`, but in addition prevents the plugin from being unloaded at runtime. If a user attempts to do so with `UNINSTALL PLUGIN`, an error occurs.

Plugin activation states are visible in the `LOAD_OPTION` column of the `INFORMATION_SCHEMA.PLUGINS` table.

Suppose that `CSV`, `BLACKHOLE`, and `ARCHIVE` are built-in pluggable storage engines and that you want the server to load them at startup, subject to these conditions: The server is permitted to run if `CSV` initialization fails, must require that `BLACKHOLE` initialization succeeds, and should disable `ARCHIVE`. To accomplish that, use these lines in an option file:

```
[mysqld]
csv=ON
blackhole=FORCE
archive=OFF
```

The `--enable-plugin_name` option format is a synonym for `--plugin_name=ON`. The `--disable-plugin_name` and `--skip-plugin_name` option formats are synonyms for `--plugin_name=OFF`.

If a plugin is disabled, either explicitly with `OFF` or implicitly because it was enabled with `ON` but failed to initialize, aspects of server operation that require the plugin will change. For example, if the plugin implements a storage engine, existing tables for the storage engine become inaccessible, and attempts to create new tables for the storage engine result in tables that use the default storage engine unless the `NO_ENGINE_SUBSTITUTION` SQL mode is enabled to cause an error to occur instead.

Disabling a plugin may require adjustment to other options. For example, if you start the server using `--skip-innodb` to disable `InnoDB`, other `innodb_xxx` options likely will need to be omitted at startup. In addition, because `InnoDB` is the default storage engine, it will not start unless you specify another available storage engine with `--default_storage_engine`. You must also set `--default_tmp_storage_engine`.

Uninstalling Plugins

At runtime, the `UNINSTALL PLUGIN` statement disables and uninstalls a plugin known to the server. The statement unloads the plugin and removes it from the `mysql.plugin` table, if it is registered there. For this reason, `UNINSTALL PLUGIN` statement requires the `DELETE` privilege for the `mysql.plugin` table. With the plugin no longer registered in the table, the server will not load the plugin automatically for subsequent restarts.

`UNINSTALL PLUGIN` can unload a plugin regardless of whether it was loaded at runtime with `INSTALL PLUGIN` or at startup with a plugin-loading option, subject to these conditions:

- It cannot unload plugins that are built in to the server. These can be identified as those that have a library name of `NULL` in the output from `INFORMATION_SCHEMA.PLUGINS` or `SHOW PLUGINS`.

- It cannot unload plugins for which the server was started with *--plugin_name*=FORCE_PLUS_PERMANENT, which prevents plugin unloading at runtime. These can be identified from the LOAD_OPTION column of the INFORMATION_SCHEMA.PLUGINS table.

To uninstall a plugin that currently is loaded at server startup with a plugin-loading option, use this procedure.

1. Remove any options related to the plugin from the my.cnf file.
2. Restart the server.
3. Plugins normally are installed using either a plugin-loading option at startup or with INSTALL PLUGIN at runtime, but not both. However, removing options for a plugin from the my.cnf file may not be sufficient to uninstall it if at some point INSTALL PLUGIN has also been used. If the plugin still appears in the output from INFORMATION_SCHEMA.PLUGINS or SHOW PLUGINS, use UNINSTALL PLUGIN to remove it from the mysql.plugin table. Then restart the server again.

5.6.3 Obtaining Server Plugin Information

There are several ways to determine which plugins are installed in the server:

- The INFORMATION_SCHEMA.PLUGINS table contains a row for each loaded plugin. Any that have a PLUGIN_LIBRARY value of NULL are built in and cannot be unloaded.

```
mysql> SELECT * FROM information_schema.PLUGINS\G
*************************** 1. row ***************************
           PLUGIN_NAME: binlog
        PLUGIN_VERSION: 1.0
         PLUGIN_STATUS: ACTIVE
           PLUGIN_TYPE: STORAGE ENGINE
   PLUGIN_TYPE_VERSION: 50158.0
        PLUGIN_LIBRARY: NULL
PLUGIN_LIBRARY_VERSION: NULL
         PLUGIN_AUTHOR: MySQL AB
    PLUGIN_DESCRIPTION: This is a pseudo storage engine to represent the binlog in a transaction
        PLUGIN_LICENSE: GPL
           LOAD_OPTION: FORCE
...
*************************** 10. row ***************************
           PLUGIN_NAME: InnoDB
        PLUGIN_VERSION: 1.0
         PLUGIN_STATUS: ACTIVE
           PLUGIN_TYPE: STORAGE ENGINE
   PLUGIN_TYPE_VERSION: 50158.0
        PLUGIN_LIBRARY: ha_innodb_plugin.so
PLUGIN_LIBRARY_VERSION: 1.0
         PLUGIN_AUTHOR: Innobase Oy
    PLUGIN_DESCRIPTION: Supports transactions, row-level locking,
                        and foreign keys
        PLUGIN_LICENSE: GPL
           LOAD_OPTION: ON
...
```

- The SHOW PLUGINS statement displays a row for each loaded plugin. Any that have a Library value of NULL are built in and cannot be unloaded.

```
mysql> SHOW PLUGINS\G
*************************** 1. row ***************************
   Name: binlog
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: NULL
License: GPL
...
*************************** 10. row ***************************
   Name: InnoDB
```

```
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: ha_innodb_plugin.so
License: GPL
...
```

- The `mysql.plugin` table shows which plugins have been registered with `INSTALL PLUGIN`. The table contains only plugin names and library file names, so it does not provide as much information as the `PLUGINS` table or the `SHOW PLUGINS` statement.

5.6.4 The Rewriter Query Rewrite Plugin

MySQL Server supports query rewrite plugins that can examine and possibly modify statements received by the server before the server executes them. See Query Rewrite Plugins.

MySQL distributions include a postparse query rewrite plugin named `Rewriter` and scripts for installing the plugin and its associated components. These components work together to provide `SELECT` rewriting capability:

- A server-side plugin named `Rewriter` examines `SELECT` statements and may rewrite them, based on its in-memory cache of rewrite rules. Standalone `SELECT` statements and `SELECT` statements in prepared statements are subject to rewriting. `SELECT` statements occurring within view definitions or stored programs are not subject to rewriting.
- The `Rewriter` plugin uses a database named `query_rewrite` containing a table named `rewrite_rules`. The table provides persistent storage for the rules that the plugin uses to decide whether to rewrite statements. Users communicate with the plugin by modifying the set of rules stored in this table. The plugin communicates with users by setting the `message` column of table rows.
- The `query_rewrite` database contains a stored procedure named `flush_rewrite_rules()` that loads the contents of the rules table into the plugin.
- A user-defined function named `load_rewrite_rules()` is used by the `flush_rewrite_rules()` stored procedure.
- The `Rewriter` plugin exposes system variables that enable plugin configuration and status variables that provide runtime operational information.

The following sections describe how to install and use the `Rewriter` plugin, and provide reference information for its associated components.

5.6.4.1 Installing or uninstalling the Rewriter Query Rewrite Plugin

Note

If installed, the `Rewriter` plugin involves some overhead even when disabled. To avoid this overhead, do not install the plugin unless you plan to use it.

To install or uninstall the `Rewriter` query rewrite plugin, choose the approropriate script located in the `share` directory of your MySQL installation:

- `install_rewriter.sql`: Choose this script to install the `Rewriter` plugin and its associated components.

Note

The `pattern_digest` and `normalized_columns` columns are also created. (For details about the table columns, see Rewriter Query Rewrite Plugin Rules Table.)

- `uninstall_rewriter.sql`: Choose this script to uninstall the `Rewriter` plugin and its associated components.

Run the chosen script as follows:

```
shell> mysql -u root -p < install_rewriter.sql
Enter password: (enter root password here)
```

The example here uses the `install_rewriter.sql` installation script. Make the appropriate substitution if you choose a different script.

Running an installation script should install and enable the plugin. To verify that, connect to the server and execute this statement:

```
mysql> SHOW GLOBAL VARIABLES LIKE 'rewriter_enabled';
+------------------+-------+
| Variable_name    | Value |
+------------------+-------+
| rewriter_enabled | ON    |
+------------------+-------+
```

For usage instructions, see Section 5.6.4.2, “Using the Rewriter Query Rewrite Plugin”. For reference information, see Section 5.6.4.3, “Rewriter Query Rewrite Plugin Reference”.

5.6.4.2 Using the Rewriter Query Rewrite Plugin

To enable or disable the plugin, enable or disable the `rewriter_enabled` system variable. By default, the `Rewriter` plugin is enabled when you install it (see Section 5.6.4.1, “Installing or uninstalling the Rewriter Query Rewrite Plugin”). To set the initial plugin state explicitly, you can set the variable at server startup. For example, to enable the plugin in an option file, use these lines:

```
[mysqld]
rewriter_enabled=ON
```

It is also possible to enable or disable the plugin at runtime:

```
mosql> SET GLOBAL rewriter_enabled = ON;
mysql> SET GLOBAL rewriter_enabled = OFF;
```

Assuming that the `Rewriter` plugin is enabled, it examines and possibly modifies each `SELECT` statement received by the server. The plugin determines whether to rewrite statements based on its in-memory cache of rewriting rules, which are loaded from the `rewrite_rules` table in the `query_rewrite` database.

Adding Rewrite Rules

To add rules for the `Rewriter` plugin, add rows to the `rewrite_rules` table, then invoke the `flush_rewrite_rules()` stored procedure to load the rules from the table into the plugin. The following example creates a simple rule to match statements that select a single literal value:

```
mysql> INSERT INTO query_rewrite.rewrite_rules (pattern, replacement)
    -> VALUES('SELECT ?', 'SELECT ? + 1');
```

The resulting table contents look like this:

```
mysql> SELECT * FROM query_rewrite.rewrite_rules\G
*************************** 1. row ***************************
                  id: 1
             pattern: SELECT ?
  pattern_database: NULL
         replacement: SELECT ? + 1
             enabled: YES
```

```
           message: NULL
    pattern_digest: NULL
normalized_pattern: NULL
```

The rule specifies a pattern template indicating which `SELECT` statements to match, and a replacement template indicating how to rewrite matching statements. However, adding the rule to the `rewrite_rules` table is not sufficient to cause the `Rewriter` plugin to use the rule. You must invoke `flush_rewrite_rules()` to load the table contents into the plugin in-memory cache:

```
mysql> CALL query_rewrite.flush_rewrite_rules();
```

Tip

If your rewrite rules seem not to be working properly, make sure that you have reloaded the rules table by calling `flush_rewrite_rules()`.

When the plugin reads each rule from the rules table, it computes a normalized form (digest) from the pattern and a digest hash value, and updates the `normalized_pattern` and `pattern_digest` columns:

```
mysql> SELECT * FROM query_rewrite.rewrite_rules\G
*************************** 1. row ***************************
                id: 1
           pattern: SELECT ?
  pattern_database: NULL
       replacement: SELECT ? + 1
           enabled: YES
           message: NULL
    pattern_digest: 46b876e64cd5c41009d91c754921f1d4
normalized_pattern: select ?
```

For information about statement digesting and normalized statements, see Section 25.9, “Performance Schema Statement Digests”.

Patterns use the same syntax as prepared statements (see Section 13.5.1, “PREPARE Syntax”). Within a pattern template, ? characters act as parameter markers that match data values. Parameter markers can be used only where data values should appear, not for SQL keywords, identifiers, and so forth. The ? characters should not be enclosed within quotation marks.

Like the pattern, the replacement can contain ? characters. For a statement that matches a pattern template, the plugin rewrites it, replacing ? parameter markers in the replacement using data values matched by the corresponding markers in the pattern. The result is a complete statement string. The plugin asks the server to parse it, and returns the result to the server as the representation of the rewritten statement.

After adding and loading the rule, check whether rewriting occurs according to whether statements match the rule pattern:

```
mysql> SELECT PI();
+----------+
| PI()     |
+----------+
| 3.141593 |
+----------+
1 row in set (0.01 sec)

mysql> SELECT 10;
+--------+
| 10 + 1 |
+--------+
|     11 |
+--------+
1 row in set, 1 warning (0.00 sec)
```

No rewriting occurs for the first `SELECT` statement, but does for the second. The second statement illustrates that when the `Rewriter` plugin rewrites a statement, it produces a warning message. To view the message, use `SHOW WARNINGS`:

```
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1105
Message: Query 'SELECT 10' rewritten to 'SELECT 10 + 1' by a query rewrite plugin
```

To enable or disable an existing rule, modify its `enabled` column and reload the table into the plugin. To disable rule 1:

```
mysql> UPDATE query_rewrite.rewrite_rules SET enabled = 'NO' WHERE id = 1;
mysql> CALL query_rewrite.flush_rewrite_rules();
```

This enables you to deactivate a rule without removing it from the table.

To re-enable rule 1:

```
mysql> UPDATE query_rewrite.rewrite_rules SET enabled = 'YES' WHERE id = 1;
mysql> CALL query_rewrite.flush_rewrite_rules();
```

The `rewrite_rules` table contains a `pattern_database` column that `Rewriter` uses for matching table names that are not qualified with a database name:

- Qualified table names in statements match qualified names in the pattern if corresponding database and table names are identical.
- Unqualified table names in statements match unqualified names in the pattern only if the default database is the same as `pattern_database` and the table names are identical.

Suppose that a table named `appdb.users` has a column named `id` and that applications are expected to select rows from the table using a query of one of these forms, where the second can be used only if `appdb` is the default database:

```
SELECT * FROM users WHERE appdb.id = id_value;
SELECT * FROM users WHERE id = id_value;
```

Suppose also that the `id` column is renamed to `user_id` (perhaps the table must be modified to add another type of ID and it is necessary to indicate more specifically what type of ID the `id` column represents).

The change means that applications must refer to `user_id` rather than `id` in the `WHERE` clause. But if there are old applications that cannot be written to change the `SELECT` queries they generate, they will no longer work properly. The `Rewriter` plugin can solve this problem. To match and rewrite statements whether or not they qualify the table name, add the following two rules and reload the rules table:

```
mysql> INSERT INTO query_rewrite.rewrite_rules
    -> (pattern, replacement) VALUES(
    -> 'SELECT * FROM appdb.users WHERE id = ?',
    -> 'SELECT * FROM appdb.users WHERE user_id = ?'
    -> );
mysql> INSERT INTO query_rewrite.rewrite_rules
    -> (pattern, replacement, pattern_database) VALUES(
    -> 'SELECT * FROM users WHERE id = ?',
    -> 'SELECT * FROM users WHERE user_id = ?',
    -> 'appdb'
    -> );
mysql> CALL query_rewrite.flush_rewrite_rules();
```

`Rewriter` uses the first rule to match statements that use the qualified table name. It uses the second to match statements that used the unqualified name, but only if the default database is `appdb` (the value in `pattern_database`).

How Statement Matching Works

The `Rewriter` plugin uses statement digests to match incoming statements against rewrite rules in stages. The `max_digest_length` system variable determines the size of the buffer used for computing statement digests. Larger values enable computation of digests that distinguish longer statements. Smaller values use less memory but increase the likelihood of longer statements colliding with the same digest value.

The plugin matches each statement to the rewrite rules as follows:

1. Compute the statement digest hash value and compare it to the rule digest hash values. This is subject to false positives, but serves as a quick rejection test.
2. If the statement digest hash value matches any pattern digest hash values, match the normalized form of the statement to the normalized form of the matching rule patterns.
3. If the normalized statement matches a rule, compare the literal values in the statement and the pattern. A `?` in the pattern matches any literal value in the statement. If the statement prepares a `SELECT` statement, `?` in the pattern also matches `?` in the statement. Otherwise, corresponding literals must be the same.

If multiple rules match a statement, it is indeterminate which one the plugin uses to rewrite the statement.

If a pattern contains more markers than the replacement, the plugin discards excess data values. If a pattern contains fewer markers than the replacement, it is an error. The plugin notices this when the rules table is loaded, writes an error message to the `message` column of the rule row to communicate the problem, and sets the `Rewriter_reload_error` status variable to `ON`.

Rewriting Prepared Statements

Prepared statements are rewritten at parse time (that is, when they are prepared), not when they are executed later.

Prepared statements differ from nonprepared statements in that they may contain `?` characters as parameter markers. To match a `?` in a prepared statement, a `Rewriter` pattern must contain `?` in the same location. Suppose that a rewrite rule has this pattern:

```
SELECT ?, 3
```

The following table shows several prepared `SELECT` statements and whether the rule pattern matches them.

Prepared Statement	Whether Pattern Matches Statement
`PREPARE s AS 'SELECT 3, 3'`	Yes
`PREPARE s AS 'SELECT ?, 3'`	Yes
`PREPARE s AS 'SELECT 3, ?'`	No
`PREPARE s AS 'SELECT ?, ?'`	No

Rewriter Plugin Operational Information

The `Rewriter` plugin makes information available about its operation by means of several status variables:

```
mysql> SHOW GLOBAL STATUS LIKE 'Rewriter%';
+----------------------------------+-------+
| Variable_name                    | Value |
+----------------------------------+-------+
| Rewriter_number_loaded_rules     | 1     |
| Rewriter_number_reloads          | 5     |
| Rewriter_number_rewritten_queries | 1     |
| Rewriter_reload_error            | ON    |
+----------------------------------+-------+
```

For descriptions of these variables, see Rewriter Query Rewrite Plugin Status Variables.

When you load the rules table by calling the `flush_rewrite_rules()` stored procedure, if an error occurs for some rule, the `CALL` statement produces an error, and the plugin sets the `Rewriter_reload_error` status variable to `ON`:

```
mysql> CALL query_rewrite.flush_rewrite_rules();
ERROR 1644 (45000): Loading of some rule(s) failed.

mysql> SHOW GLOBAL STATUS LIKE 'Rewriter_reload_error';
+-----------------------+-------+
| Variable_name         | Value |
+-----------------------+-------+
| Rewriter_reload_error | ON    |
+-----------------------+-------+
```

In this case, check the `message` column of `rewrite_rules` table rows for non-`NULL` values to see what the problem was.

Rewriter Plugin Use of Character Sets

When the `rewrite_rules` table is loaded into the `Rewriter` plugin, the plugin interprets statements using the current global value of the `character_set_client` system variable. If the global `character_set_client` value is changed subsequently, the rules table must be reloaded.

A client must have a session `character_set_client` value identical to what the global value was when the rules table was loaded or rule matching will not work for that client.

5.6.4.3 Rewriter Query Rewrite Plugin Reference

The following discussion serves as a reference to these components associated with the `Rewriter` query rewrite plugin:

- The `Rewriter` rules table in the `query_rewrite` database
- `Rewriter` procedures and functions
- `Rewriter` system and status variables

Rewriter Query Rewrite Plugin Rules Table

The `rewrite_rules` table in the `query_rewrite` database provides persistent storage for the rules that the `Rewriter` plugin uses to decide whether to rewrite statements.

Users communicate with the plugin by modifying the set of rules stored in this table. The plugin communicates with users by setting the table's `message` column.

Note

The rules table is loaded into the plugin by the `flush_rewrite_rules` stored procedure. Unless that procedure has been called following the most recent table modification, the table contents do not necessarily correspond to the set of rules the plugin is using.

The `rewrite_rules` table has these columns:

- `id`

  The rule ID. This column is the table primary key. You can use the ID to uniquely identify any rule.

- `pattern`

  The template that indicates the pattern for statements that the rule matches. Use `?` to represent parameter markers that match data values.

- `pattern_database`

  The database used to match unqualified table names in statements. Qualified table names in statements match qualified names in the pattern if corresponding database and table names are identical. Unqualified table names in statements match unqualified names in the pattern only if the default database is the same as `pattern_database` and the table names are identical.

- `replacement`

  The template that indicates how to rewrite statements matching the `pattern` column value. Use `?` to represent parameter markers that match data values. In rewritten statements, the plugin replaces `?` parameter markers in `replacement` using data values matched by the corresponding markers in `pattern`.

- `enabled`

  Whether the rule is enabled. Load operations (performed by invoking the `flush_rewrite_rules()` stored procedure) load the rule from the table into the `Rewriter` in-memory cache only if this column is `YES`.

  This column makes it possible to deactivate a rule without removing it: Set the column to a value other than `YES` and reload the table into the plugin.

- `message`

  The plugin uses this column for communicating with users. If no error occurs when the rules table is loaded into memory, the plugin sets the `message` column to `NULL`. A non-`NULL` value indicates an error and the column contents are the error message. Errors can occur under these circumstances:

  - Either the pattern or the replacement is an incorrect SQL statement that produces syntax errors.
  - The replacement contains more `?` parameter markers than the pattern.

  If a load error occurs, the plugin also sets the `Rewriter_reload_error` status variable to `ON`.

- `pattern_digest`

  This column is used for debugging and diagnostics. If the column exists when the rules table is loaded into memory, the plugin updates it with the pattern digest. This column may be useful if you are trying to determine why some statement fails to be rewritten.

- `normalized_pattern`

  This column is used for debugging and diagnostics. If the column exists when the rules table is loaded into memory, the plugin updates it with the normalized form of the pattern. This column may be useful if you are trying to determine why some statement fails to be rewritten.

Rewriter Query Rewrite Plugin Procedures and Functions

`Rewriter` plugin operation uses a stored procedure that loads the rules table into its in-memory cache, and a helper user-defined function (UDF). Under normal operation, users invoke only the stored procedure. The UDF is intended to be invoked by the stored procedure, not directly by users.

- `flush_rewrite_rules()`

  This stored procedure uses the `load_rewrite_rules()` UDF to load the contents of the `rewrite_rules` table into the `Rewriter` in-memory cache. After loading the table, it also clears the query cache.

  Calling `flush_rewrite_rules()` implies `COMMIT`.

  Invoke this procedure after you modify the rules table to cause the plugin to update its cache from the new table contents. If any errors occur, the plugin sets the `message` column for the appropriate rule rows in the table and sets the `Rewriter_reload_error` status variable to `ON`.

- `load_rewrite_rules()`

  This UDF is a helper routine used by the `flush_rewrite_rules()` stored procedure.

Rewriter Query Rewrite Plugin System Variables

The `Rewriter` query rewrite plugin supports the following system variables. These variables are available only if the plugin is installed (see Section 5.6.4.1, “Installing or uninstalling the Rewriter Query Rewrite Plugin”).

- `rewriter_enabled`

System Variable	**Name**	`rewriter_enabled`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

  Whether the `Rewriter` query rewrite plugin is enabled.

- `rewriter_verbose`

System Variable	**Name**	`rewriter_verbose`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer

  For internal use.

Rewriter Query Rewrite Plugin Status Variables

The `Rewriter` query rewrite plugin supports the following status variables. These variables are available only if the plugin is installed (see Section 5.6.4.1, “Installing or uninstalling the Rewriter Query Rewrite Plugin”).

- `Rewriter_number_loaded_rules`

  The number of rewrite plugin rewrite rules successfully loaded from the `rewrite_rules` table into memory for use by the `Rewriter` plugin.

- `Rewriter_number_reloads`

The number of times the `rewrite_rules` table has been loaded into the in-memory cache used by the `Rewriter` plugin.

- `Rewriter_number_rewritten_queries`

  The number of queries rewritten by the `Rewriter` query rewrite plugin since it was loaded.

- `Rewriter_reload_error`

  Whether an error occurred the most recent time that the `rewrite_rules` table was loaded into the in-memory cache used by the `Rewriter` plugin. If the value is `OFF`, no error occurred. If the value is `ON`, an error occurred; check the `message` column of the `rewriter_rules` table for error messages.

5.6.5 Version Tokens

MySQL includes **Version Tokens**, a feature that enables creation of and synchronization around server tokens that applications can use to prevent accessing incorrect or out-of-date data.

The Version Tokens interface has these characteristics:

- Version tokens are pairs consisting of a name that serves as a key or identifier, plus a value.
- Version tokens can be locked. An application can use token locks to indicate to other cooperating applications that tokens are in use and should not be modified.
- Version token lists are established per server; for example, to specify the server assignment or operational state. In addition, an application that communicates with a server can register its own list of tokens that indicate the state it requires the server to be in. An SQL statement sent by the application to a server not in the required state produces an error. This is a signal to the application that it should seek a different server in the required state to receive the SQL statement.

The following sections describe the components of Version Tokens, discuss how to install and use it, and provide reference information for its components.

5.6.5.1 Version Tokens Components

Version Tokens is based on a plugin library that implements these components:

- A server-side plugin named `version_tokens` holds the list of version tokens associated with the server and subscribes to notifications for statement execution events. The `version_tokens` plugin uses the audit plugin API to monitor incoming statements from clients and matches each client's session-specific version token list against the server version token list. If there is a match, the plugin lets the statement through and the server continues to process it. Otherwise, the plugin returns an error to the client and the statement fails.
- A set of user-defined functions (UDFs) provides an SQL-level API for manipulating and inspecting the list of server version tokens maintained by the plugin. The `VERSION_TOKEN_ADMIN` or `SUPER` privilege is required to call any of the Version Token UDFs.
- When the `version_tokens` plugin loads, it defines the `VERSION_TOKEN_ADMIN` dynamic privilege. This privilege can be granted to users of the UDFs.
- A system variable enables clients to specify the list of version tokens that register the required server state. If the server has a different state when a client sends a statement, the client receives an error.

5.6.5.2 Installing or Uninstalling Version Tokens

Note

If installed, Version Tokens involves some overhead. To avoid this overhead, do not install it unless you plan to use it.

This section describes how to install or uninstall Version Tokens, which is implemented in a plugin library file containing a plugin and user-defined functions. For general information about installing or uninstalling plugins and UDFs, see Section 5.6.2, “Installing and Uninstalling Plugins”, and Section 28.4.2.5, “UDF Compiling and Installing”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base name is `version_tokens`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To install the Version Tokens plugin and UDFs, use the `INSTALL PLUGIN` and `CREATE FUNCTION` statements (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN version_tokens SONAME 'version_token.so';
CREATE FUNCTION version_tokens_set RETURNS STRING SONAME 'version_token.so';
CREATE FUNCTION version_tokens_show RETURNS STRING SONAME 'version_token.so';
CREATE FUNCTION version_tokens_edit RETURNS STRING SONAME 'version_token.so';
CREATE FUNCTION version_tokens_delete RETURNS STRING SONAME 'version_token.so';
CREATE FUNCTION version_tokens_lock_shared RETURNS INT SONAME 'version_token.so';
CREATE FUNCTION version_tokens_lock_exclusive RETURNS INT SONAME 'version_token.so';
CREATE FUNCTION version_tokens_unlock RETURNS INT SONAME 'version_token.so';
```

You must install the UDFs to manage the server's version token list, but you must also install the plugin because the UDFs will not work correctly without it.

If the plugin and UDFs are used on a master replication server, install them on all slave servers as well to avoid replication problems.

Once installed as just described, the Version Tokens plugin and UDFs remain installed until uninstalled. To remove them, use the `UNINSTALL PLUGIN` and `DROP FUNCTION` statements:

```
UNINSTALL PLUGIN version_tokens;
DROP FUNCTION version_tokens_set;
DROP FUNCTION version_tokens_show;
DROP FUNCTION version_tokens_edit;
DROP FUNCTION version_tokens_delete;
DROP FUNCTION version_tokens_lock_shared;
DROP FUNCTION version_tokens_lock_exclusive;
DROP FUNCTION version_tokens_unlock;
```

5.6.5.3 Using Version Tokens

Before using Version Tokens, install it according to the instructions provided at Section 5.6.5.2, “Installing or Uninstalling Version Tokens”.

A scenario in which Version Tokens can be useful is a system that accesses a collection of MySQL servers but needs to manage them for load balancing purposes by monitoring them and adjusting server assignments according to load changes. Such a system comprises these components:

- The collection of MySQL servers to be managed.
- An administrative or management application that communicates with the servers and organizes them into high-availability groups. Groups serve different purposes, and servers within each group may have different assignments. Assignment of a server within a certain group can change at any time.
- Client applications that access the servers to retrieve and update data, choosing servers according to the purposes assigned them. For example, a client should not send an update to a read-only server.

Version Tokens permit server access to be managed according to assignment without requiring clients to repeatedly query the servers about their assignments:

- The management application performs server assignments and establishes version tokens on each server to reflect its assignment. The application caches this information to provide a central access point to it.

  If at some point the management application needs to change a server assignment (for example, to change it from permitting writes to read only), it changes the server's version token list and updates its cache.

- To improve performance, client applications obtain cache information from the management application, enabling them to avoid having to retrieve information about server assignments for each statement. Based on the type of statements it will issue (for example, reads versus writes), a client selects an appropriate server and connects to it.

- In addition, the client sends to the server its own client-specific version tokens to register the assignment it requires of the server. For each statement sent by the client to the server, the server compares its own token list with the client token list. If the server token list contains all tokens present in the client token list with the same values, there is a match and the server executes the statement.

  On the other hand, perhaps the management application has changed the server assignment and its version token list. In this case, the new server assignment may now be incompatible with the client requirements. A token mismatch between the server and client token lists occurs and the server returns an error in reply to the statement. This is an indication to the client to refresh its version token information from the management application cache, and to select a new server to communicate with.

The client-side logic for detecting version token errors and selecting a new server can be implemented different ways:

- The client can handle all version token registration, mismatch detection, and connection switching itself.

- The logic for those actions can be implemented in a connector that manages connections between clients and MySQL servers. Such a connector might handle mismatch error detection and statement resending itself, or it might pass the error to the application and leave it to the application to resend the statement.

The following example illustrates the preceding discussion in more concrete form.

When Version Tokens initializes on a given server, the server's version token list is empty. Token list maintenance is performed by calling user-defined functions (UDFs). The `VERSION_TOKEN_ADMIN` or `SUPER` privilege is required to call any of the Version Token UDFs, so token list modification is expected to be done by a management or administrative application that has that privilege.

Suppose that a management application communicates with a set of servers that are queried by clients to access employee and product databases (named `emp` and `prod`, respectively). All servers are permitted to process data retrieval statements, but only some of them are permitted to make database updates. To handle this on a database-specific basis, the management application establishes a list of version tokens on each server. In the token list for a given server, token names represent database names and token values are `read` or `write` depending on whether the database must be used in read-only fashion or whether it can take reads and writes.

Client applications register a list of version tokens they require the server to match by setting a system variable. Variable setting occurs on a client-specific basis, so different clients can register different requirements. By default, the client token list is empty, which matches any server token list. When a client sets its token list to a nonempty value, matching may succeed or fail, depending on the server version token list.

To define the version token list for a server, the management application calls the `version_token_set()` UDF. (There are also UDFs for modifying and displaying the token list, described later.) For example, the application might send these statements to a group of three servers:

Server 1:

```
mysql> SELECT version_tokens_set('emp=read;prod=read');
+-----------------------------------------+
| version_tokens_set('emp=read;prod=read') |
+-----------------------------------------+
| 2 version tokens set.                   |
+-----------------------------------------+
```

Server 2:

```
mysql> SELECT version_tokens_set('emp=write;prod=read');
+------------------------------------------+
| version_tokens_set('emp=write;prod=read') |
+------------------------------------------+
| 2 version tokens set.                    |
+------------------------------------------+
```

Server 3:

```
mysql> SELECT version_tokens_set('emp=read;prod=write');
+------------------------------------------+
| version_tokens_set('emp=read;prod=write') |
+------------------------------------------+
| 2 version tokens set.                    |
+------------------------------------------+
```

The token list in each case is specified as a semicolon-separated list of *`name=value`* pairs. The resulting token list values result in these server assingments:

- Any server accepts reads for either database.
- Only server 2 accepts updates for the `emp` database.
- Only server 3 accepts updates for the `prod` database.

In addition to assigning each server a version token list, the management application also maintains a cache that reflects the server assignments.

Before communicating with the servers, a client application contacts the management application and retrieves information about server assignments. Then the client selects a server based on those assignments. Suppose that a client wants to perform both reads and writes on the `emp` database. Based on the preceding assignments, only server 2 qualifies. The client connects to server 2 and registers its server requirements there by setting its `version_tokens_session` system variable:

```
mysql> SET @@session.version_tokens_session = 'emp=write';
```

For subsequent statements sent by the client to server 2, the server compares its own version token list to the client list to check whether they match. If so, statements execute normally:

```
mysql> UPDATE emp.employee SET salary = salary * 1.1 WHERE id = 4981;
Query OK, 1 row affected (0.07 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT last_name, first_name FROM emp.employee WHERE id = 4981;
+-----------+------------+
| last_name | first_name |
```

```
+-----------+------------+
| Smith     | Abe        |
+-----------+------------+
1 row in set (0.01 sec)
```

Discrepancies between the server and client version token lists can occur two ways:

- A token name in the `version_tokens_session` value is not present in the server token list. In this case, an `ER_VTOKEN_PLUGIN_TOKEN_NOT_FOUND` error occurs.
- A token value in the `version_tokens_session` value differs from the value of the corresponding token in the server token list. In this case, an `ER_VTOKEN_PLUGIN_TOKEN_MISMATCH` error occurs.

As long as the assignment of server 2 does not change, the client continues to use it for reads and writes. But suppose that the management application wants to change server assignments so that writes for the `emp` database must be sent to server 1 instead of server 2. To do this, it uses `version_tokens_edit()` to modify the `emp` token value on the two servers (and updates its cache of server assignments):

Server 1:

```
mysql> SELECT version_tokens_edit('emp=write');
+-----------------------------------+
| version_tokens_edit('emp=write') |
+-----------------------------------+
| 1 version tokens updated.         |
+-----------------------------------+
```

Server 2:

```
mysql> SELECT version_tokens_edit('emp=read');
+--------------------------------+
| version_tokens_edit('emp=read') |
+--------------------------------+
| 1 version tokens updated.      |
+--------------------------------+
```

`version_tokens_edit()` modifies the named tokens in the server token list and leaves other tokens unchanged.

The next time the client sends a statement to server 2, its own token list no longer matches the server token list and an error occurs:

```
mysql> UPDATE emp.employee SET salary = salary * 1.1 WHERE id = 4982;
ERROR 3136 (42000): Version token mismatch for emp. Correct value read
```

In this case, the client should contact the management application to obtain updated information about server assignments, select a new server, and send the failed statement to the new server.

Note

Each client must cooperate with Version Tokens by sending only statements in accordance with the token list that it registers with a given server. For example, if a client registers a token list of `'emp=read'`, there is nothing in Version Tokens to prevent the client from sending updates for the `emp` database. The client itself must refrain from doing so.

For each statement received from a client, the server implicitly uses locking, as follows:

- Take a shared lock for each token named in the client token list (that is, in the `version_tokens_session` value)

- Perform the comparison between the server and client token lists
- Execute the statement or produce an error depending on the comparison result
- Release the locks

The server uses shared locks so that comparisons for multiple sessions can occur without blocking, while preventing changes to the tokens for any session that attempts to acquire an exclusive lock before it manipulates tokens of the same names in the server token list.

The preceding example uses only a few of the user-defined included in the Version Tokens plugin library, but there are others. One set of UDFs permits the server's list of version tokens to be manipulated and inspected. Another set of UDFs permits version tokens to be locked and unlocked.

These UDFs permit the server's list of version tokens to be created, changed, removed, and inspected:

- `version_tokens_set()` completely replaces the current list and assigns a new list. The argument is a semicolon-separated list of *`name=value`* pairs.
- `version_tokens_edit()` enables partial modifications to the current list. It can add new tokens or change the values of existing tokens. The argument is a semicolon-separated list of *`name=value`* pairs.
- `version_tokens_delete()` deletes tokens from the current list. The argument is a semicolon-separated list of token names.
- `version_tokens_show()` displays the current token list. It takes no argument.

Each of those functions, if successful, returns a binary string indicating what action occurred. The following example establishes the server token list, modifies it by adding a new token, deletes some tokens, and displays the resulting token list:

```
mysql> SELECT version_tokens_set('tok1=a;tok2=b');
+-----------------------------------+
| version_tokens_set('tok1=a;tok2=b') |
+-----------------------------------+
| 2 version tokens set.               |
+-----------------------------------+
mysql> SELECT version_tokens_edit('tok3=c');
+-------------------------------+
| version_tokens_edit('tok3=c') |
+-------------------------------+
| 1 version tokens updated.     |
+-------------------------------+
mysql> SELECT version_tokens_delete('tok2;tok1');
+-----------------------------------+
| version_tokens_delete('tok2;tok1') |
+-----------------------------------+
| 2 version tokens deleted.         |
+-----------------------------------+
mysql> SELECT version_tokens_show();
+-----------------------+
| version_tokens_show() |
+-----------------------+
| tok3=c;               |
+-----------------------+
```

Warnings occur if a token list is malformed:

```
mysql> SELECT version_tokens_set('tok1=a; =c');
+---------------------------------+
| version_tokens_set('tok1=a; =c') |
+---------------------------------+
| 1 version tokens set.           |
```

```
+---------------------------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 42000
Message: Invalid version token pair encountered. The list provided
         is only partially updated.
1 row in set (0.00 sec)
```

As mentioned previously, version tokens are defined using a semicolon-separated list of *`name=value`* pairs. Consider this invocation of `version_tokens_set()`:

```
mysql> SELECT version_tokens_set('tok1=b;;; tok2= a = b ; tok1 = 1\'2 3"4')
+-----------------------------------------------------------------+
| version_tokens_set('tok1=b;;; tok2= a = b ; tok1 = 1\'2 3"4') |
+-----------------------------------------------------------------+
| 3 version tokens set.                                           |
+-----------------------------------------------------------------+
```

Version Tokens interprets the argument as follows:

- Whitespace around names and values is ignored. Whitespace within names and values is permitted. (For `version_tokens_delete()`, which takes a list of names without values, whitespace around names is ignored.)
- There is no quoting mechanism.
- Order of tokens is not significant except that if a token list contains multiple instances of a given token name, the last value takes precedence over earlier values.

Given those rules, the preceding `version_tokens_set()` call results in a token list with two tokens: `tok1` has the value `1'2 3"4`, and `tok2` has the value `a = b`. To verify this, call `version_tokens_show()`:

```
mysql> SELECT version_tokens_show();
+--------------------------+
| version_tokens_show()    |
+--------------------------+
| tok2=a = b;tok1=1'2 3"4; |
+--------------------------+
```

If the token list contains two tokens, why did `version_tokens_set()` return the value `3 version tokens set`? That occurred because the original token list contained two definitions for `tok1`, and the second definition replaced the first.

The Version Tokens token-manipulation UDFs place these constraints on token names and values:

- Token names cannot contain `=` or `;` characters and have a maximum length of 64 characters.
- Token values cannot contain `;` characters. Length of values is constrained by the value of the `max_allowed_packet` system variable.
- Version Tokens treats token names and values as binary strings, so comparisons are case sensitive.

Version Tokens also includes a set of UDFs enabling tokens to be locked and unlocked:

- `version_tokens_lock_exclusive()` acquires exclusive version token locks. It takes a list of one or more lock names and a timeout value.
- `version_tokens_lock_shared()` acquires shared version token locks. It takes a list of one or more lock names and a timeout value.

- `version_tokens_unlock()` releases version token locks (exclusive and shared). It takes no argument.

Each locking function returns nonzero for success. Otherwise, an error occurs:

```
mysql> SELECT version_tokens_lock_shared('lock1', 'lock2', 0);
+-----------------------------------------------+
| version_tokens_lock_shared('lock1', 'lock2', 0) |
+-----------------------------------------------+
|                                             1 |
+-----------------------------------------------+

mysql> SELECT version_tokens_lock_shared(NULL, 0);
ERROR 3131 (42000): Incorrect locking service lock name '(null)'.
```

Locking using Version Tokens locking functions is advisory; applications must agree to cooperate.

It is possible to lock nonexisting token names. This does not create the tokens.

Note

Version Tokens locking functions are based on the locking service described at Section 28.3.1, “The Locking Service”, and thus have the same semantics for shared and exclusive locks. (Version Tokens uses the locking service routines built into the server, not the locking service UDF interface, so those UDFs need not be installed to use Version Tokens.) Locks acquired by Version Tokens use a locking service namespace of `version_token_locks`. Locking service locks can be monitored using the Performance Schema, so this is also true for Version Tokens locks. For details, see Locking Service Monitoring.

For the Version Tokens locking functions, token name arguments are used exactly as specified. Surrounding whitespace is not ignored and `=` and `;` characters are permitted. This is because Version Tokens simply passes the token names to be locked as is to the locking service.

5.6.5.4 Version Tokens Reference

The following discussion serves as a reference to these Version Tokens components:

- Version Tokens Functions
- Version Tokens System Variables

Version Tokens Functions

The Version Tokens plugin library includes several user-defined functions. One set of UDFs permits the server's list of version tokens to be manipulated and inspected. Another set of UDFs permits version tokens to be locked and unlocked. The `VERSION_TOKEN_ADMIN` or `SUPER` privilege is required to invoke any Version Tokens UDF.

The following UDFs permit the server's list of version tokens to be created, changed, removed, and inspected. Interpretation of *`name_list`* and *`token_list`* arguments (including whitespace handling) occurs as described in Section 5.6.5.3, “Using Version Tokens”, which provides details about the syntax for specifying tokens, as well as additional examples.

- `version_tokens_delete(`*`name_list`*`)`

  Deletes tokens from the server's list of version tokens using the *`name_list`* argument and returns a binary string that indicates the outcome of the operation. *`name_list`* is a semicolon-separated list of version token names to delete.

```
mysql> SELECT version_tokens_delete('tok1;tok3');
+-----------------------------------+
| version_tokens_delete('tok1;tok3') |
+-----------------------------------+
| 2 version tokens deleted.          |
+-----------------------------------+
```

An argument of `NULL` is treated as an empty string, which has no effect on the token list.

`version_tokens_delete()` deletes the tokens named in its argument, if they exist. (It is not an error to delete nonexisting tokens.) To clear the token list entirely without knowing which tokens are in the list, pass `NULL` or a string containing no tokens to `version_tokens_set()`:

```
mysql> SELECT version_tokens_set(NULL);
+----------------------------+
| version_tokens_set(NULL)     |
+----------------------------+
| Version tokens list cleared. |
+----------------------------+
mysql> SELECT version_tokens_set('');
+----------------------------+
| version_tokens_set('')       |
+----------------------------+
| Version tokens list cleared. |
+----------------------------+
```

- `version_tokens_edit(`*`token_list`*`)`

  Modifies the server's list of version tokens using the *`token_list`* argument and returns a binary string that indicates the outcome of the operation. *`token_list`* is a semicolon-separated list of *`name=value`* pairs specifying the name of each token to be defined and its value. If a token exists, its value is updated with the given value. If a token does not exist, it is created with the given value. If the argument is `NULL` or a string containing no tokens, the token list remains unchanged.

```
mysql> SELECT version_tokens_set('tok1=value1;tok2=value2');
+-----------------------------------------------+
| version_tokens_set('tok1=value1;tok2=value2') |
+-----------------------------------------------+
| 2 version tokens set.                         |
+-----------------------------------------------+
mysql> SELECT version_tokens_edit('tok2=new_value2;tok3=new_value3');
+-------------------------------------------------------+
| version_tokens_edit('tok2=new_value2;tok3=new_value3') |
+-------------------------------------------------------+
| 2 version tokens updated.                             |
+-------------------------------------------------------+
```

- `version_tokens_set(`*`token_list`*`)`

  Replaces the server's list of version tokens with the tokens defined in the *`token_list`* argument and returns a binary string that indicates the outcome of the operation. *`token_list`* is a semicolon-separated list of *`name=value`* pairs specifying the name of each token to be defined and its value. If the argument is `NULL` or a string containing no tokens, the token list is cleared.

```
mysql> SELECT version_tokens_set('tok1=value1;tok2=value2');
+-----------------------------------------------+
| version_tokens_set('tok1=value1;tok2=value2') |
+-----------------------------------------------+
| 2 version tokens set.                         |
+-----------------------------------------------+
```

- `version_tokens_show()`

  Returns the server's list of version tokens as a binary string containing a semicolon-separated list of *`name=value`* pairs.

```
mysql> SELECT version_tokens_show();
+-------------------------+
| version_tokens_show()   |
+-------------------------+
| tok2=value2;tok1=value1; |
+-------------------------+
```

The following UDFs permit version tokens to be locked and unlocked:

- `version_tokens_lock_exclusive(token_name[, token_name] ..., timeout)`

  Acquires exclusive locks on one or more version tokens, specified by name as strings, timing out with an error if the locks are not acquired within the given timeout value.

```
mysql> SELECT version_tokens_lock_exclusive('lock1', 'lock2', 10);
+-----------------------------------------------------+
| version_tokens_lock_exclusive('lock1', 'lock2', 10) |
+-----------------------------------------------------+
|                                                   1 |
+-----------------------------------------------------+
```

- `version_tokens_lock_shared(token_name[, token_name] ..., timeout)`

  Acquires shared locks on one or more version tokens, specified by name as strings, timing out with an error if the locks are not acquired within the given timeout value.

```
mysql> SELECT version_tokens_lock_shared('lock1', 'lock2', 10);
+-------------------------------------------------+
| version_tokens_lock_shared('lock1', 'lock2', 10) |
+-------------------------------------------------+
|                                               1 |
+-------------------------------------------------+
```

- `version_tokens_unlock()`

  Releases all locks that were acquired within the current session using `version_tokens_lock_exclusive()` and `version_tokens_lock_shared()`.

```
mysql> SELECT version_tokens_unlock();
+-------------------------+
| version_tokens_unlock() |
+-------------------------+
|                       1 |
+-------------------------+
```

The locking functions share these characteristics:

- The return value is nonzero for success. Otherwise, an error occurs.
- Token names are strings.
- In contrast to argument handling for the UDFs that manipulate the server token list, whitespace surrounding token name arguments is not ignored and `=` and `;` characters are permitted.
- It is possible to lock nonexisting token names. This does not create the tokens.
- Timeout values are nonnegative integers representing the time in seconds to wait to acquire locks before timing out with an error. If the timeout is 0, there is no waiting and the function produces an error if locks cannot be acquired immediately.
- Version Tokens locking functions are based on the locking service described at Section 28.3.1, “The Locking Service”.

Version Tokens System Variables

Version Tokens supports the following system variables. These variables are unavailable unless the Version Tokens plugin is installed (see Section 5.6.5.2, “Installing or Uninstalling Version Tokens”).

System variables:

- `version_tokens_session`

Command-Line Format	`--version-tokens-session=value`	
System Variable	Name	`version_tokens_session`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	string
	Default	`NULL`

The session value of this variable specifies the client version token list and indicates the tokens that the client session requires the server version token list to have.

If the `version_tokens_session` variable is `NULL` (the default) or has an empty value, any server version token list matches. (In effect, an empty value disables matching requirements.)

If the `version_tokens_session` variable has a nonempty value, any mismatch between its value and the server version token list results in an error for any statement the session sends to the server. A mismatch occurs under these conditions:

- A token name in the `version_tokens_session` value is not present in the server token list. In this case, an `ER_VTOKEN_PLUGIN_TOKEN_NOT_FOUND` error occurs.
- A token value in the `version_tokens_session` value differs from the value of the corresponding token in the server token list. In this case, an `ER_VTOKEN_PLUGIN_TOKEN_MISMATCH` error occurs.

It is not a mismatch for the server version token list to include a token not named in the `version_tokens_session` value.

Suppose that a management application has set the server token list as follows:

```
mysql> SELECT version_tokens_set('tok1=a;tok2=b;tok3=c');
+-----------------------------------------+
| version_tokens_set('tok1=a;tok2=b;tok3=c') |
+-----------------------------------------+
| 3 version tokens set.                   |
+-----------------------------------------+
```

A client registers the tokens it requires the server to match by setting its `version_tokens_session` value. Then, for each subsequent statement sent by the client, the server checks its token list against the client `version_tokens_session` value and produces an error if there is a mismatch:

```
mysql> SET @@session.version_tokens_session = 'tok1=a;tok2=b';
mysql> SELECT 1;
+---+
| 1 |
+---+
| 1 |
+---+
```

```
mysql> SET @@session.version_tokens_session = 'tok1=b';
mysql> SELECT 1;
ERROR 3136 (42000): Version token mismatch for tok1. Correct value a
```

The first `SELECT` succeeds because the client tokens `tok1` and `tok2` are present in the server token list and each token has the same value in the server list. The second `SELECT` fails because, although `tok1` is present in the server token list, it has a different value than specified by the client.

At this point, any statement sent by the client fails, unless the server token list changes such that it matches again. Suppose that the management application changes the server token list as follows:

```
mysql> SELECT version_tokens_edit('tok1=b');
+------------------------------+
| version_tokens_edit('tok1=b') |
+------------------------------+
| 1 version tokens updated.     |
+------------------------------+
mysql> SELECT version_tokens_show();
+-----------------------+
| version_tokens_show() |
+-----------------------+
| tok3=c;tok1=b;tok2=b; |
+-----------------------+
```

Now the client `version_tokens_session` value matches the server token list and the client can once again successfully execute statements:

```
mysql> SELECT 1;
+---+
| 1 |
+---+
| 1 |
+---+
```

- `version_tokens_session_number`

Command-Line Format	`--version-tokens-session-number=N`	
System Variable	Name	`version_tokens_session_number`
	Variable Scope	Global, Session
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`0`

This variable is for internal use.

5.7 Running Multiple MySQL Instances on One Machine

In some cases, you might want to run multiple instances of MySQL on a single machine. You might want to test a new MySQL release while leaving an existing production setup undisturbed. Or you might want to give different users access to different `mysqld` servers that they manage themselves. (For example, you might be an Internet Service Provider that wants to provide independent MySQL installations for different customers.)

It is possible to use a different MySQL server binary per instance, or use the same binary for multiple instances, or any combination of the two approaches. For example, you might run a server from MySQL 5.7 and one from MySQL 8.0, to see how different versions handle a given workload. Or

you might run multiple instances of the current production version, each managing a different set of databases.

Whether or not you use distinct server binaries, each instance that you run must be configured with unique values for several operating parameters. This eliminates the potential for conflict between instances. Parameters can be set on the command line, in option files, or by setting environment variables. See Section 4.2.3, “Specifying Program Options”. To see the values used by a given instance, connect to it and execute a `SHOW VARIABLES` statement.

The primary resource managed by a MySQL instance is the data directory. Each instance should use a different data directory, the location of which is specified using the `--datadir=dir_name` option. For methods of configuring each instance with its own data directory, and warnings about the dangers of failing to do so, see Section 5.7.1, “Setting Up Multiple Data Directories”.

In addition to using different data directories, several other options must have different values for each server instance:

- `--port=port_num`

  `--port` controls the port number for TCP/IP connections. Alternatively, if the host has multiple network addresses, you can use `--bind-address` to cause each server to listen to a different address.

- `--socket={file_name|pipe_name}`

  `--socket` controls the Unix socket file path on Unix or the named pipe name on Windows. On Windows, it is necessary to specify distinct pipe names only for those servers configured to permit named-pipe connections.

- `--shared-memory-base-name=name`

  This option is used only on Windows. It designates the shared-memory name used by a Windows server to permit clients to connect using shared memory. It is necessary to specify distinct shared-memory names only for those servers configured to permit shared-memory connections.

- `--pid-file=file_name`

  This option indicates the path name of the file in which the server writes its process ID.

If you use the following log file options, their values must differ for each server:

- `--general_log_file=file_name`
- `--log-bin[=file_name]`
- `--slow_query_log_file=file_name`
- `--log-error[=file_name]`

For further discussion of log file options, see Section 5.4, “MySQL Server Logs”.

To achieve better performance, you can specify the following option differently for each server, to spread the load between several physical disks:

- `--tmpdir=dir_name`

Having different temporary directories also makes it easier to determine which MySQL server created any given temporary file.

If you have multiple MySQL installations in different locations, you can specify the base directory for each installation with the `--basedir=dir_name` option. This causes each instance to automatically use a different data directory, log files, and PID file because the default for each of those parameters

is relative to the base directory. In that case, the only other options you need to specify are the `--socket` and `--port` options. Suppose that you install different versions of MySQL using `tar` file binary distributions. These install in different locations, so you can start the server for each installation using the command `bin/mysqld_safe` under its corresponding base directory. `mysqld_safe` determines the proper `--basedir` option to pass to `mysqld`, and you need specify only the `--socket` and `--port` options to `mysqld_safe`.

As discussed in the following sections, it is possible to start additional servers by specifying appropriate command options or by setting environment variables. However, if you need to run multiple servers on a more permanent basis, it is more convenient to use option files to specify for each server those option values that must be unique to it. The `--defaults-file` option is useful for this purpose.

5.7.1 Setting Up Multiple Data Directories

Each MySQL Instance on a machine should have its own data directory. The location is specified using the `--datadir=dir_name` option.

There are different methods of setting up a data directory for a new instance:

- Create a new data directory.
- Copy an existing data directory.

The following discussion provides more detail about each method.

Warning

Normally, you should never have two servers that update data in the same databases. This may lead to unpleasant surprises if your operating system does not support fault-free system locking. If (despite this warning) you run multiple servers using the same data directory and they have logging enabled, you must use the appropriate options to specify log file names that are unique to each server. Otherwise, the servers try to log to the same files.

Even when the preceding precautions are observed, this kind of setup works only with `MyISAM` and `MERGE` tables, and not with any of the other storage engines. Also, this warning against sharing a data directory among servers always applies in an NFS environment. Permitting multiple MySQL servers to access a common data directory over NFS is a *very bad idea*. The primary problem is that NFS is the speed bottleneck. It is not meant for such use. Another risk with NFS is that you must devise a way to ensure that two or more servers do not interfere with each other. Usually NFS file locking is handled by the `lockd` daemon, but at the moment there is no platform that performs locking 100% reliably in every situation.

Create a New Data Directory

With this method, the data directory will be in the same state as when you first install MySQL. It will have the default set of MySQL accounts and no user data.

On Unix, initialize the data directory. See Section 2.9, “Postinstallation Setup and Testing”.

On Windows, the data directory is included in the MySQL distribution:

- MySQL Zip archive distributions for Windows contain an unmodified data directory. You can unpack such a distribution into a temporary location, then copy it `data` directory to where you are setting up the new instance.
- Windows MSI package installers create and set up the data directory that the installed server will use, but also create a pristine “template” data directory named `data` under the installation directory.

After an installation has been performed using an MSI package, the template data directory can be copied to set up additional MySQL instances.

Copy an Existing Data Directory

With this method, any MySQL accounts or user data present in the data directory are carried over to the new data directory.

1. Stop the existing MySQL instance using the data directory. This must be a clean shutdown so that the instance flushes any pending changes to disk.
2. Copy the data directory to the location where the new data directory should be.
3. Copy the `my.cnf` or `my.ini` option file used by the existing instance. This serves as a basis for the new instance.
4. Modify the new option file so that any pathnames referring to the original data directory refer to the new data directory. Also, modify any other options that must be unique per instance, such as the TCP/IP port number and the log files. For a list of parameters that must be unique per instance, see Section 5.7, “Running Multiple MySQL Instances on One Machine”.
5. Start the new instance, telling it to use the new option file.

5.7.2 Running Multiple MySQL Instances on Windows

You can run multiple servers on Windows by starting them manually from the command line, each with appropriate operating parameters, or by installing several servers as Windows services and running them that way. General instructions for running MySQL from the command line or as a service are given in Section 2.3, “Installing MySQL on Microsoft Windows”. The following sections describe how to start each server with different values for those options that must be unique per server, such as the data directory. These options are listed in Section 5.7, “Running Multiple MySQL Instances on One Machine”.

5.7.2.1 Starting Multiple MySQL Instances at the Windows Command Line

The procedure for starting a single MySQL server manually from the command line is described in Section 2.3.5.6, “Starting MySQL from the Windows Command Line”. To start multiple servers this way, you can specify the appropriate options on the command line or in an option file. It is more convenient to place the options in an option file, but it is necessary to make sure that each server gets its own set of options. To do this, create an option file for each server and tell the server the file name with a `--defaults-file` option when you run it.

Suppose that you want to run one instance of `mysqld` on port 3307 with a data directory of `C:\mydata1`, and another instance on port 3308 with a data directory of `C:\mydata2`. Use this procedure:

1. Make sure that each data directory exists, including its own copy of the `mysql` database that contains the grant tables.
2. Create two option files. For example, create one file named `C:\my-opts1.cnf` that looks like this:

```
[mysqld]
datadir = C:/mydata1
port = 3307
```

Create a second file named `C:\my-opts2.cnf` that looks like this:

```
[mysqld]
datadir = C:/mydata2
port = 3308
```

3. Use the `--defaults-file` option to start each server with its own option file:

```
C:\> C:\mysql\bin\mysqld --defaults-file=C:\my-opts1.cnf
C:\> C:\mysql\bin\mysqld --defaults-file=C:\my-opts2.cnf
```

   Each server starts in the foreground (no new prompt appears until the server exits later), so you will need to issue those two commands in separate console windows.

To shut down the servers, connect to each using the appropriate port number:

```
C:\> C:\mysql\bin\mysqladmin --port=3307 --host=127.0.0.1 --user=root --password shutdown
C:\> C:\mysql\bin\mysqladmin --port=3308 --host=127.0.0.1 --user=root --password shutdown
```

Servers configured as just described permit clients to connect over TCP/IP. If your version of Windows supports named pipes and you also want to permit named-pipe connections, specify options that enable the named pipe and specify its name. Each server that supports named-pipe connections must use a unique pipe name. For example, the `C:\my-opts1.cnf` file might be written like this:

```
[mysqld]
datadir = C:/mydata1
port = 3307
enable-named-pipe
socket = mypipe1
```

Modify `C:\my-opts2.cnf` similarly for use by the second server. Then start the servers as described previously.

A similar procedure applies for servers that you want to permit shared-memory connections. Enable such connections with the `--shared-memory` option and specify a unique shared-memory name for each server with the `--shared-memory-base-name` option.

5.7.2.2 Starting Multiple MySQL Instances as Windows Services

On Windows, a MySQL server can run as a Windows service. The procedures for installing, controlling, and removing a single MySQL service are described in Section 2.3.5.8, “Starting MySQL as a Windows Service”.

To set up multiple MySQL services, you must make sure that each instance uses a different service name in addition to the other parameters that must be unique per instance.

For the following instructions, suppose that you want to run the `mysqld` server from two different versions of MySQL that are installed at `C:\mysql-5.5.9` and `C:\mysql-8.0.3`, respectively. (This might be the case if you are running 5.5.9 as your production server, but also want to conduct tests using 8.0.3.)

To install MySQL as a Windows service, use the `--install` or `--install-manual` option. For information about these options, see Section 2.3.5.8, “Starting MySQL as a Windows Service”.

Based on the preceding information, you have several ways to set up multiple services. The following instructions describe some examples. Before trying any of them, shut down and remove any existing MySQL services.

- **Approach 1:** Specify the options for all services in one of the standard option files. To do this, use a different service name for each server. Suppose that you want to run the 5.5.9 `mysqld` using the service name of `mysqld1` and the 8.0.3 `mysqld` using the service name `mysqld2`. In this case, you can use the `[mysqld1]` group for 5.5.9 and the `[mysqld2]` group for 8.0.3. For example, you can set up `C:\my.cnf` like this:

```
# options for mysqld1 service
[mysqld1]
```

```
basedir = C:/mysql-5.5.9
port = 3307
enable-named-pipe
socket = mypipe1

# options for mysqld2 service
[mysqld2]
basedir = C:/mysql-8.0.3
port = 3308
enable-named-pipe
socket = mypipe2
```

Install the services as follows, using the full server path names to ensure that Windows registers the correct executable program for each service:

```
C:\> C:\mysql-5.5.9\bin\mysqld --install mysqld1
C:\> C:\mysql-8.0.3\bin\mysqld --install mysqld2
```

To start the services, use the services manager, or use `NET START` with the appropriate service names:

```
C:\> NET START mysqld1
C:\> NET START mysqld2
```

To stop the services, use the services manager, or use `NET STOP` with the appropriate service names:

```
C:\> NET STOP mysqld1
C:\> NET STOP mysqld2
```

- **Approach 2:** Specify options for each server in separate files and use `--defaults-file` when you install the services to tell each server what file to use. In this case, each file should list options using a `[mysqld]` group.

  With this approach, to specify options for the 5.5.9 `mysqld`, create a file `C:\my-opts1.cnf` that looks like this:

  ```
  [mysqld]
  basedir = C:/mysql-5.5.9
  port = 3307
  enable-named-pipe
  socket = mypipe1
  ```

  For the 8.0.3 `mysqld`, create a file `C:\my-opts2.cnf` that looks like this:

  ```
  [mysqld]
  basedir = C:/mysql-8.0.3
  port = 3308
  enable-named-pipe
  socket = mypipe2
  ```

  Install the services as follows (enter each command on a single line):

  ```
  C:\> C:\mysql-5.5.9\bin\mysqld --install mysqld1
            --defaults-file=C:\my-opts1.cnf
  C:\> C:\mysql-8.0.3\bin\mysqld --install mysqld2
            --defaults-file=C:\my-opts2.cnf
  ```

  When you install a MySQL server as a service and use a `--defaults-file` option, the service name must precede the option.

  After installing the services, start and stop them the same way as in the preceding example.

To remove multiple services, use `mysqld --remove` for each one, specifying a service name following the `--remove` option. If the service name is the default (`MySQL`), you can omit it.

5.7.3 Running Multiple MySQL Instances on Unix

Note

The discussion here uses `mysqld_safe` to launch multiple instances of MySQL. For MySQL installation using an RPM distribution, server startup and shutdown is managed by systemd on several Linux platforms. On these platforms, `mysqld_safe` is not installed because it is unnecessary. For information about using systemd to handle multiple MySQL instances, see Section 2.5.9, “Managing MySQL Server with systemd”.

One way is to run multiple MySQL instances on Unix is to compile different servers with different default TCP/IP ports and Unix socket files so that each one listens on different network interfaces. Compiling in different base directories for each installation also results automatically in a separate, compiled-in data directory, log file, and PID file location for each server.

Assume that an existing 5.7 server is configured for the default TCP/IP port number (3306) and Unix socket file (`/tmp/mysql.sock`). To configure a new 8.0.3 server to have different operating parameters, use a `CMake` command something like this:

```
shell> cmake . -DMYSQL_TCP_PORT=port_number \
             -DMYSQL_UNIX_ADDR=file_name \
             -DCMAKE_INSTALL_PREFIX=/usr/local/mysql-8.0.3
```

Here, *`port_number`* and *`file_name`* must be different from the default TCP/IP port number and Unix socket file path name, and the `CMAKE_INSTALL_PREFIX` value should specify an installation directory different from the one under which the existing MySQL installation is located.

If you have a MySQL server listening on a given port number, you can use the following command to find out what operating parameters it is using for several important configurable variables, including the base directory and Unix socket file name:

```
shell> mysqladmin --host=host_name --port=port_number variables
```

With the information displayed by that command, you can tell what option values *not* to use when configuring an additional server.

If you specify `localhost` as the host name, `mysqladmin` defaults to using a Unix socket file connection rather than TCP/IP. To explicitly specify the connection protocol, use the `--protocol={TCP|SOCKET|PIPE|MEMORY}` option.

You need not compile a new MySQL server just to start with a different Unix socket file and TCP/IP port number. It is also possible to use the same server binary and start each invocation of it with different parameter values at runtime. One way to do so is by using command-line options:

```
shell> mysqld_safe --socket=file_name --port=port_number
```

To start a second server, provide different `--socket` and `--port` option values, and pass a `--datadir=dir_name` option to `mysqld_safe` so that the server uses a different data directory.

Alternatively, put the options for each server in a different option file, then start each server using a `--defaults-file` option that specifies the path to the appropriate option file. For example, if the option files for two server instances are named `/usr/local/mysql/my.cnf` and `/usr/local/mysql/my.cnf2`, start the servers like this: command:

```
shell> mysqld_safe --defaults-file=/usr/local/mysql/my.cnf
```

```
shell> mysqld_safe --defaults-file=/usr/local/mysql/my.cnf2
```

Another way to achieve a similar effect is to use environment variables to set the Unix socket file name and TCP/IP port number:

```
shell> MYSQL_UNIX_PORT=/tmp/mysqld-new.sock
shell> MYSQL_TCP_PORT=3307
shell> export MYSQL_UNIX_PORT MYSQL_TCP_PORT
shell> bin/mysqld --initialize --user=mysql
shell> mysqld_safe --datadir=/path/to/datadir &
```

This is a quick way of starting a second server to use for testing. The nice thing about this method is that the environment variable settings apply to any client programs that you invoke from the same shell. Thus, connections for those clients are automatically directed to the second server.

Section 4.9, “MySQL Program Environment Variables”, includes a list of other environment variables you can use to affect MySQL programs.

On Unix, the `mysqld_multi` script provides another way to start multiple servers. See Section 4.3.4, “`mysqld_multi` — Manage Multiple MySQL Servers”.

5.7.4 Using Client Programs in a Multiple-Server Environment

To connect with a client program to a MySQL server that is listening to different network interfaces from those compiled into your client, you can use one of the following methods:

- Start the client with `--host=host_name --port=port_number` to connect using TCP/IP to a remote server, with `--host=127.0.0.1 --port=port_number` to connect using TCP/IP to a local server, or with `--host=localhost --socket=file_name` to connect to a local server using a Unix socket file or a Windows named pipe.
- Start the client with `--protocol=TCP` to connect using TCP/IP, `--protocol=SOCKET` to connect using a Unix socket file, `--protocol=PIPE` to connect using a named pipe, or `--protocol=MEMORY` to connect using shared memory. For TCP/IP connections, you may also need to specify `--host` and `--port` options. For the other types of connections, you may need to specify a `--socket` option to specify a Unix socket file or Windows named-pipe name, or a `--shared-memory-base-name` option to specify the shared-memory name. Shared-memory connections are supported only on Windows.
- On Unix, set the `MYSQL_UNIX_PORT` and `MYSQL_TCP_PORT` environment variables to point to the Unix socket file and TCP/IP port number before you start your clients. If you normally use a specific socket file or port number, you can place commands to set these environment variables in your `.login` file so that they apply each time you log in. See Section 4.9, “MySQL Program Environment Variables”.
- Specify the default Unix socket file and TCP/IP port number in the `[client]` group of an option file. For example, you can use `C:\my.cnf` on Windows, or the `.my.cnf` file in your home directory on Unix. See Section 4.2.6, “Using Option Files”.
- In a C program, you can specify the socket file or port number arguments in the `mysql_real_connect()` call. You can also have the program read option files by calling `mysql_options()`. See Section 27.7.7, “C API Function Descriptions”.
- If you are using the Perl `DBD::mysql` module, you can read options from MySQL option files. For example:

```
$dsn = "DBI:mysql:test;mysql_read_default_group=client;"
       . "mysql_read_default_file=/usr/local/mysql/data/my.cnf";
$dbh = DBI->connect($dsn, $user, $password);
```

See Section 27.9, “MySQL Perl API”.

Other programming interfaces may provide similar capabilities for reading option files.

Chapter 6 Security

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When thinking about security within a MySQL installation, you should consider a wide range of possible topics and how they affect the security of your MySQL server and related applications:

- General factors that affect security. These include choosing good passwords, not granting unnecessary privileges to users, ensuring application security by preventing SQL injections and data corruption, and others. See Section 6.1, “General Security Issues”.

- Security of the installation itself. The data files, log files, and the all the application files of your installation should be protected to ensure that they are not readable or writable by unauthorized parties. For more information, see Section 2.9, “Postinstallation Setup and Testing”.
- Access control and security within the database system itself, including the users and databases granted with access to the databases, views and stored programs in use within the database. For more information, see Section 6.2, “The MySQL Access Privilege System”, and Section 6.3, “MySQL User Account Management”.
- The features offered by security-related plugins. See Section 6.5, “Security Plugins”.
- Network security of MySQL and your system. The security is related to the grants for individual users, but you may also wish to restrict MySQL so that it is available only locally on the MySQL server host, or to a limited set of other hosts.
- Ensure that you have adequate and appropriate backups of your database files, configuration and log files. Also be sure that you have a recovery solution in place and test that you are able to successfully recover the information from your backups. See Chapter 7, *Backup and Recovery*.

6.1 General Security Issues

This section describes general security issues to be aware of and what you can do to make your MySQL installation more secure against attack or misuse. For information specifically about the access control system that MySQL uses for setting up user accounts and checking database access, see Section 2.9, “Postinstallation Setup and Testing”.

For answers to some questions that are often asked about MySQL Server security issues, see Section A.9, “MySQL 8.0 FAQ: Security”.

6.1.1 Security Guidelines

Anyone using MySQL on a computer connected to the Internet should read this section to avoid the most common security mistakes.

In discussing security, it is necessary to consider fully protecting the entire server host (not just the MySQL server) against all types of applicable attacks: eavesdropping, altering, playback, and denial of service. We do not cover all aspects of availability and fault tolerance here.

MySQL uses security based on Access Control Lists (ACLs) for all connections, queries, and other operations that users can attempt to perform. There is also support for SSL-encrypted connections between MySQL clients and servers. Many of the concepts discussed here are not specific to MySQL at all; the same general ideas apply to almost all applications.

When running MySQL, follow these guidelines:

- **Do not ever give anyone (except MySQL `root` accounts) access to the `user` table in the `mysql` database!** This is critical.
- Learn how the MySQL access privilege system works (see Section 6.2, “The MySQL Access Privilege System”). Use the `GRANT` and `REVOKE` statements to control access to MySQL. Do not grant more privileges than necessary. Never grant privileges to all hosts.

  Checklist:

  - Try `mysql -u root`. If you are able to connect successfully to the server without being asked for a password, anyone can connect to your MySQL server as the MySQL `root` user with full privileges! Review the MySQL installation instructions, paying particular attention to the information about setting a `root` password. See Section 2.9.4, “Securing the Initial MySQL Account”.
  - Use the `SHOW GRANTS` statement to check which accounts have access to what. Then use the `REVOKE` statement to remove those privileges that are not necessary.

- Do not store cleartext passwords in your database. If your computer becomes compromised, the intruder can take the full list of passwords and use them. Instead, use `SHA2()` or some other one-way hashing function and store the hash value.

  To prevent password recovery using rainbow tables, do not use these functions on a plain password; instead, choose some string to be used as a salt, and use hash(hash(password)+salt) values.

- Do not choose passwords from dictionaries. Special programs exist to break passwords. Even passwords like “xfish98” are very bad. Much better is “duag98” which contains the same word “fish” but typed one key to the left on a standard QWERTY keyboard. Another method is to use a password that is taken from the first characters of each word in a sentence (for example, “Four score and seven years ago” results in a password of “Fsasya”). The password is easy to remember and type, but difficult to guess for someone who does not know the sentence. In this case, you can additionally substitute digits for the number words to obtain the phrase “4 score and 7 years ago”, yielding the password “4sa7ya” which is even more difficult to guess.

- Invest in a firewall. This protects you from at least 50% of all types of exploits in any software. Put MySQL behind the firewall or in a demilitarized zone (DMZ).

  Checklist:

  - Try to scan your ports from the Internet using a tool such as `nmap`. MySQL uses port 3306 by default. This port should not be accessible from untrusted hosts. As a simple way to check whether your MySQL port is open, try the following command from some remote machine, where *`server_host`* is the host name or IP address of the host on which your MySQL server runs:

    ```
    shell> telnet server_host 3306
    ```

    If `telnet` hangs or the connection is refused, the port is blocked, which is how you want it to be. If you get a connection and some garbage characters, the port is open, and should be closed on your firewall or router, unless you really have a good reason to keep it open.

- Applications that access MySQL should not trust any data entered by users, and should be written using proper defensive programming techniques. See Section 6.1.7, “Client Programming Security Guidelines”.

- Do not transmit plain (unencrypted) data over the Internet. This information is accessible to everyone who has the time and ability to intercept it and use it for their own purposes. Instead, use an encrypted protocol such as SSL or SSH. MySQL supports internal SSL connections. Another technique is to use SSH port-forwarding to create an encrypted (and compressed) tunnel for the communication.

- Learn to use the `tcpdump` and `strings` utilities. In most cases, you can check whether MySQL data streams are unencrypted by issuing a command like the following:

  ```
  shell> tcpdump -l -i eth0 -w - src or dst port 3306 | strings
  ```

  This works under Linux and should work with small modifications under other systems.

Warning

If you do not see cleartext data, this does not always mean that the information actually is encrypted. If you need high security, consult with a security expert.

6.1.2 Keeping Passwords Secure

Passwords occur in several contexts within MySQL. The following sections provide guidelines that enable end users and administrators to keep these passwords secure and avoid exposing them. In

addition, the `validate_password` plugin can be used to enforce a policy on acceptable password. See Section 6.5.3, “The Password Validation Plugin”.

6.1.2.1 End-User Guidelines for Password Security

MySQL users should use the following guidelines to keep passwords secure.

When you run a client program to connect to the MySQL server, it is inadvisable to specify your password in a way that exposes it to discovery by other users. The methods you can use to specify your password when you run client programs are listed here, along with an assessment of the risks of each method. In short, the safest methods are to have the client program prompt for the password or to specify the password in a properly protected option file.

- Use the `mysql_config_editor` utility, which enables you to store authentication credentials in an encrypted login path file named `.mylogin.cnf`. The file can be read later by MySQL client programs to obtain authentication credentials for connecting to MySQL Server. See Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

- Use a `-p`*`your_pass`* or `--password=`*`your_pass`* option on the command line. For example:

```
shell> mysql -u francis -pfrank db_name
```

  This is convenient *but insecure*. On some systems, your password becomes visible to system status programs such as `ps` that may be invoked by other users to display command lines. MySQL clients typically overwrite the command-line password argument with zeros during their initialization sequence. However, there is still a brief interval during which the value is visible. Also, on some systems this overwriting strategy is ineffective and the password remains visible to `ps`. (SystemV Unix systems and perhaps others are subject to this problem.)

  If your operating environment is set up to display your current command in the title bar of your terminal window, the password remains visible as long as the command is running, even if the command has scrolled out of view in the window content area.

- Use the `-p` or `--password` option on the command line with no password value specified. In this case, the client program solicits the password interactively:

```
shell> mysql -u francis -p db_name
Enter password: ********
```

  The `*` characters indicate where you enter your password. The password is not displayed as you enter it.

  It is more secure to enter your password this way than to specify it on the command line because it is not visible to other users. However, this method of entering a password is suitable only for programs that you run interactively. If you want to invoke a client from a script that runs noninteractively, there is no opportunity to enter the password from the keyboard. On some systems, you may even find that the first line of your script is read and interpreted (incorrectly) as your password.

- Store your password in an option file. For example, on Unix, you can list your password in the `[client]` section of the `.my.cnf` file in your home directory:

```
[client]
password=your_pass
```

  To keep the password safe, the file should not be accessible to anyone but yourself. To ensure this, set the file access mode to `400` or `600`. For example:

```
shell> chmod 600 .my.cnf
```

To name from the command line a specific option file containing the password, use the `--defaults-file=file_name` option, where `file_name` is the full path name to the file. For example:

```
shell> mysql --defaults-file=/home/francis/mysql-opts
```

Section 4.2.6, “Using Option Files”, discusses option files in more detail.

- Store your password in the `MYSQL_PWD` environment variable. See Section 4.9, “MySQL Program Environment Variables”.

  This method of specifying your MySQL password must be considered *extremely insecure* and should not be used. Some versions of `ps` include an option to display the environment of running processes. On some systems, if you set `MYSQL_PWD`, your password is exposed to any other user who runs `ps`. Even on systems without such a version of `ps`, it is unwise to assume that there are no other methods by which users can examine process environments.

On Unix, the `mysql` client writes a record of executed statements to a history file (see Section 4.5.1.3, “mysql Logging”). By default, this file is named `.mysql_history` and is created in your home directory. Passwords can be written as plain text in SQL statements such as `CREATE USER` and `ALTER USER`, so if you use these statements, they are logged in the history file. To keep this file safe, use a restrictive access mode, the same way as described earlier for the `.my.cnf` file.

If your command interpreter is configured to maintain a history, any file in which the commands are saved will contain MySQL passwords entered on the command line. For example, `bash` uses `~/.bash_history`. Any such file should have a restrictive access mode.

6.1.2.2 Administrator Guidelines for Password Security

Database administrators should use the following guidelines to keep passwords secure.

MySQL stores passwords for user accounts in the `mysql.user` table. Access to this table should never be granted to any nonadministrative accounts.

Account passwords can be expired so that users must reset them. See Section 6.3.8, “Password Expiration Policy”, and Section 6.3.9, “Password Expiration and Sandbox Mode”.

The `validate_password` plugin can be used to enforce a policy on acceptable password. See Section 6.5.3, “The Password Validation Plugin”.

A user who has access to modify the plugin directory (the value of the `plugin_dir` system variable) or the `my.cnf` file that specifies the plugin directory location can replace plugins and modify the capabilities provided by plugins, including authentication plugins.

Files such as log files to which passwords might be written should be protected. See Section 6.1.2.3, “Passwords and Logging”.

6.1.2.3 Passwords and Logging

Passwords can be written as plain text in SQL statements such as `CREATE USER`, `GRANT`, `SET PASSWORD`, and statements that invoke the `PASSWORD()` function. If such statements are logged by the MySQL server as written, passwords in them become visible to anyone with access to the logs.

In MySQL 8.0, statement logging avoids writing passwords in cleartext for the following statements:

```
CREATE USER ... IDENTIFIED BY ...
ALTER USER ... IDENTIFIED BY ...
GRANT ... IDENTIFIED BY ...
SET PASSWORD ...
SLAVE START ... PASSWORD = ...
CREATE SERVER ... OPTIONS(... PASSWORD ...)
```

```
ALTER SERVER ... OPTIONS(... PASSWORD ...)
```

Passwords in those statements are rewritten to not appear literally in statement text written to the general query log, slow query log, and binary log. Rewriting does not apply to other statements. In particular, `INSERT` or `UPDATE` statements for the `mysql.user` table that refer to literal passwords are logged as is, so you should avoid such statements. (Direct manipulation of grant tables is discouraged, anyway.)

For the general query log, password rewriting can be suppressed by starting the server with the `--log-raw` option. For security reasons, this option is not recommended for production use. For diagnostic purposes, it may be useful to see the exact text of statements as received by the server.

Statements received by the server may be rewritten if a query rewrite plugin is installed (see Query Rewrite Plugins). In this case, the `--log-raw` option affects statement logging as follows:

- Without `--log-raw`, the server logs the statement returned by the query rewrite plugin. This may differ from the statement as received.
- With `--log-raw`, the server logs the original statement as received.

An implication of password rewriting is that statements that cannot be parsed (due, for example, to syntax errors) are not written to the general query log because they cannot be known to be password free. Use cases that require logging of all statements including those with errors should use the `--log-raw` option, bearing in mind that this also bypasses password rewriting.

Password rewriting occurs only when plain text passwords are expected. For statements with syntax that expect a password hash value, no rewriting occurs. If a plain text password is supplied erroneously for such syntax, the password is logged as given, without rewriting. For example, the following statement is logged as shown because a password hash value is expected:

```
CREATE USER 'user1'@'localhost' IDENTIFIED BY PASSWORD 'not-so-secret';
```

To guard log files against unwarranted exposure, locate them in a directory that restricts access to the server and the database administrator. If the server logs to tables in the `mysql` database, grant access to those tables only to the database administrator.

Replication slaves store the password for the replication master in the master info repository, which can be either a file or a table (see Section 18.2.4, “Replication Relay and Status Logs”). Ensure that the repository can be accessed only by the database administrator. An alternative to storing the password in a file is to use the `START SLAVE` statement to specify credentials for connecting to the master.

Use a restricted access mode to protect database backups that include log tables or log files containing passwords.

6.1.3 Making MySQL Secure Against Attackers

When you connect to a MySQL server, you should use a password. The password is not transmitted in clear text over the connection. Password handling during the client connection sequence was upgraded in MySQL 4.1.1 to be very secure. If you are still using pre-4.1.1-style passwords, the encryption algorithm is not as strong as the newer algorithm. With some effort, a clever attacker who can sniff the traffic between the client and the server can crack the password.

All other information is transferred as text, and can be read by anyone who is able to watch the connection. If the connection between the client and the server goes through an untrusted network, and you are concerned about this, you can use the compressed protocol to make traffic much more difficult to decipher. You can also use MySQL's internal SSL support to make the connection even more secure. See Section 6.4, “Using Secure Connections”. Alternatively, use SSH to get an encrypted TCP/IP connection between a MySQL server and a MySQL client. You can find an Open Source SSH client at http://www.openssh.org/, and a comparison of both Open Source and Commercial SSH clients at http://en.wikipedia.org/wiki/Comparison_of_SSH_clients.

To make a MySQL system secure, you should strongly consider the following suggestions:

- Require all MySQL accounts to have a password. A client program does not necessarily know the identity of the person running it. It is common for client/server applications that the user can specify any user name to the client program. For example, anyone can use the `mysql` program to connect as any other person simply by invoking it as `mysql -u other_user db_name` if `other_user` has no password. If all accounts have a password, connecting using another user's account becomes much more difficult.

  For a discussion of methods for setting passwords, see Section 6.3.7, “Assigning Account Passwords”.

- Make sure that the only Unix user account with read or write privileges in the database directories is the account that is used for running `mysqld`.

- Never run the MySQL server as the Unix `root` user. This is extremely dangerous, because any user with the `FILE` privilege is able to cause the server to create files as `root` (for example, `~root/.bashrc`). To prevent this, `mysqld` refuses to run as `root` unless that is specified explicitly using the `--user=root` option.

  `mysqld` can (and should) be run as an ordinary, unprivileged user instead. You can create a separate Unix account named `mysql` to make everything even more secure. Use this account only for administering MySQL. To start `mysqld` as a different Unix user, add a `user` option that specifies the user name in the `[mysqld]` group of the `my.cnf` option file where you specify server options. For example:

  ```
  [mysqld]
  user=mysql
  ```

  This causes the server to start as the designated user whether you start it manually or by using `mysqld_safe` or `mysql.server`. For more details, see Section 6.1.5, “How to Run MySQL as a Normal User”.

  Running `mysqld` as a Unix user other than `root` does not mean that you need to change the `root` user name in the `user` table. *User names for MySQL accounts have nothing to do with user names for Unix accounts*.

- Do not grant the `FILE` privilege to nonadministrative users. Any user that has this privilege can write a file anywhere in the file system with the privileges of the `mysqld` daemon. This includes the server's data directory containing the files that implement the privilege tables. To make `FILE`-privilege operations a bit safer, files generated with `SELECT ... INTO OUTFILE` do not overwrite existing files and are writable by everyone.

  The `FILE` privilege may also be used to read any file that is world-readable or accessible to the Unix user that the server runs as. With this privilege, you can read any file into a database table. This could be abused, for example, by using `LOAD DATA` to load `/etc/passwd` into a table, which then can be displayed with `SELECT`.

  To limit the location in which files can be read and written, set the `secure_file_priv` system to a specific directory. See Section 5.1.5, “Server System Variables”.

- Do not grant the `PROCESS` or `SUPER` privilege to nonadministrative users. The output of `mysqladmin processlist` and `SHOW PROCESSLIST` shows the text of any statements currently being executed, so any user who is permitted to see the server process list might be able to see statements issued by other users such as `UPDATE user SET password=PASSWORD('not_secure')`.

  `mysqld` reserves an extra connection for users who have the `SUPER` privilege, so that a MySQL `root` user can log in and check server activity even if all normal connections are in use.

  The `SUPER` privilege can be used to terminate client connections, change server operation by changing the value of system variables, and control replication servers.

- Do not permit the use of symlinks to tables. (This capability can be disabled with the `--skip-symbolic-links` option.) This is especially important if you run `mysqld` as `root`, because anyone that has write access to the server's data directory then could delete any file in the system! See Section 8.12.2.2, “Using Symbolic Links for MyISAM Tables on Unix”.
- Stored programs and views should be written using the security guidelines discussed in Section 23.6, “Access Control for Stored Programs and Views”.
- If you do not trust your DNS, you should use IP addresses rather than host names in the grant tables. In any case, you should be very careful about creating grant table entries using host name values that contain wildcards.
- If you want to restrict the number of connections permitted to a single account, you can do so by setting the `max_user_connections` variable in `mysqld`. The `GRANT` statement also supports resource control options for limiting the extent of server use permitted to an account. See Section 13.7.1.6, “GRANT Syntax”.
- If the plugin directory is writable by the server, it may be possible for a user to write executable code to a file in the directory using `SELECT ... INTO DUMPFILE`. This can be prevented by making `plugin_dir` read only to the server or by setting `--secure-file-priv` to a directory where `SELECT` writes can be made safely.

6.1.4 Security-Related mysqld Options and Variables

The following table shows `mysqld` options and system variables that affect security. For descriptions of each of these, see Section 5.1.4, “Server Command Options”, and Section 5.1.5, “Server System Variables”.

Table 6.1 Security Option/Variable Summary

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
allow-suspicious-udfs	Yes	Yes				
automatic_sp_privileges			Yes		Global	Yes
chroot	Yes	Yes				
des-key-file	Yes	Yes				
local_infile			Yes		Global	Yes
old_passwords			Yes		Both	Yes
safe-user-create	Yes	Yes				
secure-auth	Yes	Yes			Global	Yes
- *Variable*: secure_auth			Yes		Global	Yes
secure-file-priv	Yes	Yes			Global	No
- *Variable*: secure_file_priv			Yes		Global	No
skip-grant-tables	Yes	Yes				
skip-name-resolve	Yes	Yes			Global	No
- *Variable*: skip_name_resolve			Yes		Global	No
skip-networking	Yes	Yes			Global	No
- *Variable*: skip_networking			Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
skip-show-database	Yes	Yes			Global	No
- *Variable*: skip_show_database			Yes		Global	No

6.1.5 How to Run MySQL as a Normal User

On Windows, you can run the server as a Windows service using a normal user account.

On Linux, for installations performed using a MySQL repository or RPM packages, the MySQL server `mysqld` should be started by the local `mysql` operating system user. Starting by another operating system user is not supported by the init scripts that are included as part of the MySQL repositories.

On Unix (or Linux for installations performed using `tar.gz` packages) , the MySQL server `mysqld` can be started and run by any user. However, you should avoid running the server as the Unix `root` user for security reasons. To change `mysqld` to run as a normal unprivileged Unix user *`user_name`*, you must do the following:

1. Stop the server if it is running (use `mysqladmin shutdown`).
2. Change the database directories and files so that *`user_name`* has privileges to read and write files in them (you might need to do this as the Unix `root` user):

   ```
   shell> chown -R user_name /path/to/mysql/datadir
   ```

   If you do not do this, the server will not be able to access databases or tables when it runs as *`user_name`*.

   If directories or files within the MySQL data directory are symbolic links, `chown -R` might not follow symbolic links for you. If it does not, you will also need to follow those links and change the directories and files they point to.
3. Start the server as user *`user_name`*. Another alternative is to start `mysqld` as the Unix `root` user and use the `--user=`*`user_name`* option. `mysqld` starts up, then switches to run as the Unix user *`user_name`* before accepting any connections.
4. To start the server as the given user automatically at system startup time, specify the user name by adding a `user` option to the `[mysqld]` group of the `/etc/my.cnf` option file or the `my.cnf` option file in the server's data directory. For example:

   ```
   [mysqld]
   user=user_name
   ```

If your Unix machine itself is not secured, you should assign passwords to the MySQL `root` account in the grant tables. Otherwise, any user with a login account on that machine can run the `mysql` client with a `--user=root` option and perform any operation. (It is a good idea to assign passwords to MySQL accounts in any case, but especially so when other login accounts exist on the server host.) See Section 2.9.4, “Securing the Initial MySQL Account”.

6.1.6 Security Issues with LOAD DATA LOCAL

The `LOAD DATA` statement can load a file located on the server host, or, if the `LOCAL` keyword is specified, on the client host.

There are two potential security issues with the `LOCAL` version of `LOAD DATA`:

- The transfer of the file from the client host to the server host is initiated by the MySQL server. In theory, a patched server could be built that would tell the client program to transfer a file of the server's choosing rather than the file named by the client in the `LOAD DATA` statement. Such a

server could access any file on the client host to which the client user has read access. (A patched server could in fact reply with a file-transfer request to any statement, not just `LOAD DATA LOCAL`, so a more fundamental issue is that clients should not connect to untrusted servers.)

- In a Web environment where the clients are connecting from a Web server, a user could use `LOAD DATA LOCAL` to read any files that the Web server process has read access to (assuming that a user could run any statement against the SQL server). In this environment, the client with respect to the MySQL server actually is the Web server, not a remote program being run by users who connect to the Web server.

To avoid `LOAD DATA` issues, clients should avoid using `LOCAL`. To avoid connecting to untrusted servers, clients can establish a secure connection and verify the server identity by connecting using the `--ssl-mode=VERIFY_IDENTIFY` option and the appropriate CA certificate.

To enable adminstrators and applications to manage the local data loading capability, `LOCAL` configuration works like this:

- On the server side:
  - The `local_infile` system variable controls server-side `LOCAL` capability. Depending on the `local_infile` setting, the server refuses or permits local data loading by clients that have `LOCAL` enabled on the client side. By default, `local_infile` is disabled. (Enabled prior to MySQL 8.0.2.)
  - To explicitly cause the server to refuse or permit `LOAD DATA LOCAL` statements (regardless of how client programs and libraries are configured at build time or runtime), start `mysqld` with `local_infile` disabled or enabled, respectively. `local_infile` can also be set at runtime.
- On the client side:
  - The `ENABLED_LOCAL_INFILE` `CMake` option controls the compiled-in default `LOCAL` capability for the MySQL client library. Clients that make no explicit arrangements therefore have `LOCAL` capability disabled or enabled according to the `ENABLED_LOCAL_INFILE` setting specified at MySQL build time.

    By default, the client library in MySQL binary distributions is compiled with `ENABLED_LOCAL_INFILE` disabled. (Enabled prior to MySQL 8.0.2.) If you compile MySQL from source, configure it with `ENABLED_LOCAL_INFILE` disabled or enabled based on whether clients that make no explicit arrangements should have `LOCAL` capability disabled or enabled, respectively.
  - Client programs that use the C API can control load data loading explicitly by invoking `mysql_options()` to disable or enable the `MYSQL_OPT_LOCAL_INFILE` option. See Section 27.7.7.50, “mysql_options()”.
  - For the `mysql` client, local data loading is disabled by default. To disable or enable it explicitly, use the `--local-infile=0` or `--local-infile[=1]` option.
  - For the `mysqlimport` client, local data loading is disabled by default. To disable or enable it explicitly, use the `--local=0` or `--local[=1]` option.
  - If you use `LOAD DATA LOCAL` in Perl scripts or other programs that read the `[client]` group from option files, you can add an `local-infile` option setting to that group. To prevent problems for programs that do not understand this option, specify it using the `loose-` prefix:

    ```
    [client]
    loose-local-infile=0
    ```

    or:

    ```
    ```

```
[client]
loose-local-infile=1
```

- In all cases, successful use of a `LOCAL` load operation by a client also requires that the server permits it.

If `LOCAL` capability is disabled, on either the server or client side, a client that attempts to issue a `LOAD DATA LOCAL` statement receives the following error message:

```
ERROR 1148: The used command is not allowed with this MySQL version
```

6.1.7 Client Programming Security Guidelines

Applications that access MySQL should not trust any data entered by users, who can try to trick your code by entering special or escaped character sequences in Web forms, URLs, or whatever application you have built. Be sure that your application remains secure if a user enters something like `; DROP DATABASE mysql;`. This is an extreme example, but large security leaks and data loss might occur as a result of hackers using similar techniques, if you do not prepare for them.

A common mistake is to protect only string data values. Remember to check numeric data as well. If an application generates a query such as `SELECT * FROM table WHERE ID=234` when a user enters the value `234`, the user can enter the value `234 OR 1=1` to cause the application to generate the query `SELECT * FROM table WHERE ID=234 OR 1=1`. As a result, the server retrieves every row in the table. This exposes every row and causes excessive server load. The simplest way to protect from this type of attack is to use single quotation marks around the numeric constants: `SELECT * FROM table WHERE ID='234'`. If the user enters extra information, it all becomes part of the string. In a numeric context, MySQL automatically converts this string to a number and strips any trailing nonnumeric characters from it.

Sometimes people think that if a database contains only publicly available data, it need not be protected. This is incorrect. Even if it is permissible to display any row in the database, you should still protect against denial of service attacks (for example, those that are based on the technique in the preceding paragraph that causes the server to waste resources). Otherwise, your server becomes unresponsive to legitimate users.

Checklist:

- Enable strict SQL mode to tell the server to be more restrictive of what data values it accepts. See Section 5.1.8, “Server SQL Modes”.
- Try to enter single and double quotation marks (`'` and `"`) in all of your Web forms. If you get any kind of MySQL error, investigate the problem right away.
- Try to modify dynamic URLs by adding `%22` (`"`), `%23` (`#`), and `%27` (`'`) to them.
- Try to modify data types in dynamic URLs from numeric to character types using the characters shown in the previous examples. Your application should be safe against these and similar attacks.
- Try to enter characters, spaces, and special symbols rather than numbers in numeric fields. Your application should remove them before passing them to MySQL or else generate an error. Passing unchecked values to MySQL is very dangerous!
- Check the size of data before passing it to MySQL.
- Have your application connect to the database using a user name different from the one you use for administrative purposes. Do not give your applications any access privileges they do not need.

Many application programming interfaces provide a means of escaping special characters in data values. Properly used, this prevents application users from entering values that cause the application to generate statements that have a different effect than you intend:

- MySQL C API: Use the `mysql_real_escape_string_quote()` API call.

- MySQL++: Use the `escape` and `quote` modifiers for query streams.
- PHP: Use either the `mysqli` or `pdo_mysql` extensions, and not the older `ext/mysql` extension. The preferred API's support the improved MySQL authentication protocol and passwords, as well as prepared statements with placeholders. See also Choosing an API.

  If the older `ext/mysql` extension must be used, then for escaping use the `mysql_real_escape_string_quote()` function and not `mysql_escape_string()` or `addslashes()` because only `mysql_real_escape_string_quote()` is character set-aware; the other functions can be “bypassed” when using (invalid) multibyte character sets.
- Perl DBI: Use placeholders or the `quote()` method.
- Ruby DBI: Use placeholders or the `quote()` method.
- Java JDBC: Use a `PreparedStatement` object and placeholders.

Other programming interfaces might have similar capabilities.

6.2 The MySQL Access Privilege System

The primary function of the MySQL privilege system is to authenticate a user who connects from a given host and to associate that user with privileges on a database such as `SELECT`, `INSERT`, `UPDATE`, and `DELETE`. Additional functionality includes the ability to have anonymous users and to grant privileges for MySQL-specific functions such as `LOAD DATA INFILE` and administrative operations.

There are some things that you cannot do with the MySQL privilege system:

- You cannot explicitly specify that a given user should be denied access. That is, you cannot explicitly match a user and then refuse the connection.
- You cannot specify that a user has privileges to create or drop tables in a database but not to create or drop the database itself.
- A password applies globally to an account. You cannot associate a password with a specific object such as a database, table, or routine.

The user interface to the MySQL privilege system consists of SQL statements such as `CREATE USER`, `GRANT`, and `REVOKE`. See Section 13.7.1, “Account Management Statements”.

Internally, the server stores privilege information in the grant tables of the `mysql` database (that is, in the database named `mysql`). The MySQL server reads the contents of these tables into memory when it starts and bases access-control decisions on the in-memory copies of the grant tables.

The MySQL privilege system ensures that all users may perform only the operations permitted to them. As a user, when you connect to a MySQL server, your identity is determined by *the host from which you connect* and *the user name you specify*. When you issue requests after connecting, the system grants privileges according to your identity and *what you want to do*.

MySQL considers both your host name and user name in identifying you because there is no reason to assume that a given user name belongs to the same person on all hosts. For example, the user `joe` who connects from `office.example.com` need not be the same person as the user `joe` who connects from `home.example.com`. MySQL handles this by enabling you to distinguish users on different hosts that happen to have the same name: You can grant one set of privileges for connections by `joe` from `office.example.com`, and a different set of privileges for connections by `joe` from `home.example.com`. To see what privileges a given account has, use the `SHOW GRANTS` statement. For example:

```
SHOW GRANTS FOR 'joe'@'office.example.com';
SHOW GRANTS FOR 'joe'@'home.example.com';
```

MySQL access control involves two stages when you run a client program that connects to the server:

Stage 1: The server accepts or rejects the connection based on your identity and whether you can verify your identity by supplying the correct password.

Stage 2: Assuming that you can connect, the server checks each statement you issue to determine whether you have sufficient privileges to perform it. For example, if you try to select rows from a table in a database or drop a table from the database, the server verifies that you have the `SELECT` privilege for the table or the `DROP` privilege for the database.

For a more detailed description of what happens during each stage, see Section 6.2.6, “Access Control, Stage 1: Connection Verification”, and Section 6.2.7, “Access Control, Stage 2: Request Verification”.

If your privileges are changed (either by yourself or someone else) while you are connected, those changes do not necessarily take effect immediately for the next statement that you issue. For details about the conditions under which the server reloads the grant tables, see Section 6.2.8, “When Privilege Changes Take Effect”.

For general security-related advice, see Section 6.1, “General Security Issues”. For help in diagnosing privilege-related problems, see Section 6.2.9, “Troubleshooting Problems Connecting to MySQL”.

6.2.1 Privileges Provided by MySQL

The privileges granted to a MySQL account determine which operations the account can perform. MySQL privileges differ in the contexts in which they apply and at different levels of operation:

- Administrative privileges enable users to manage operation of the MySQL server. These privileges are global because they are not specific to a particular database.
- Database privileges apply to a database and to all objects within it. These privileges can be granted for specific databases, or globally so that they apply to all databases.
- Privileges for database objects such as tables, indexes, views, and stored routines can be granted for specific objects within a database, for all objects of a given type within a database (for example, all tables in a database), or globally for all objects of a given type in all databases).

Privileges also differ in terms of whether they are static (built in to the server) or dynamic (defined at runtime). Whether a privilege is static or dynamic affects its availability to be granted to user accounts and roles. See Section 6.2.2, “Static Versus Dynamic Privileges”.

Information about account privileges is stored in the `user`, `db`, `tables_priv`, `columns_priv`, `procs_priv`, and `global_grants` tables in the `mysql` system database (see Section 6.2.3, “Grant Tables”). The MySQL server reads the contents of these tables into memory when it starts and reloads them under the circumstances indicated in Section 6.2.8, “When Privilege Changes Take Effect”. Access-control decisions are based on the in-memory copies of the grant tables.

Some MySQL releases introduce changes to the structure of the grant tables to add new privileges or features. To make sure that you can take advantage of any new capabilities, update your grant tables to have the current structure whenever you upgrade MySQL. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

The following tables show the static and dynamic privilege names used in `GRANT` and `REVOKE` statements, along with the column name associated with each privilege in the grant tables and the context in which the privilege applies.

Table 6.2 Permissible Static Privileges for GRANT and REVOKE

Privilege	Column	Context
`ALL [PRIVILEGES]`	Synonym for “all privileges”	Server administration
`ALTER`	`Alter_priv`	Tables
`ALTER ROUTINE`	`Alter_routine_priv`	Stored routines

Privilege	Column	Context
CREATE	Create_priv	Databases, tables, or indexes
CREATE ROLE	Create_role_priv	Server administration
CREATE ROUTINE	Create_routine_priv	Stored routines
CREATE TABLESPACE	Create_tablespace_priv	Server administration
CREATE TEMPORARY TABLES	Create_tmp_table_priv	Tables
CREATE USER	Create_user_priv	Server administration
CREATE VIEW	Create_view_priv	Views
DELETE	Delete_priv	Tables
DROP	Drop_priv	Databases, tables, or views
DROP ROLE	Drop_role_priv	Server administration
EVENT	Event_priv	Databases
EXECUTE	Execute_priv	Stored routines
FILE	File_priv	File access on server host
GRANT OPTION	Grant_priv	Databases, tables, or stored routines
INDEX	Index_priv	Tables
INSERT	Insert_priv	Tables or columns
LOCK TABLES	Lock_tables_priv	Databases
PROCESS	Process_priv	Server administration
PROXY	See proxies_priv table	Server administration
REFERENCES	References_priv	Databases or tables
RELOAD	Reload_priv	Server administration
REPLICATION CLIENT	Repl_client_priv	Server administration
REPLICATION SLAVE	Repl_slave_priv	Server administration
SELECT	Select_priv	Tables or columns
SHOW DATABASES	Show_db_priv	Server administration
SHOW VIEW	Show_view_priv	Views
SHUTDOWN	Shutdown_priv	Server administration
SUPER	Super_priv	Server administration
TRIGGER	Trigger_priv	Tables
UPDATE	Update_priv	Tables or columns
USAGE	Synonym for “no privileges”	Server administration

Table 6.3 Permissible Dynamic Privileges for GRANT and REVOKE

Privilege	Column	Context
AUDIT_ADMIN	See global_grants table	Server administration
BINLOG_ADMIN	See global_grants table	Server administration
CONNECTION_ADMIN	See global_grants table	Server administration
ENCRYPTION_KEY_ADMIN	See global_grants table	Server administration
FIREWALL_ADMIN	See global_grants table	Server administration
FIREWALL_USER	See global_grants table	Server administration
GROUP_REPLICATION_ADMIN	See global_grants table	Server administration

Privilege	Column	Context
REPLICATION_SLAVE_ADMIN	See global_grants table	Server administration
ROLE_ADMIN	See global_grants table	Server administration
SET_USER_ID	See global_grants table	Server administration
SYSTEM_VARIABLES_ADMIN	See global_grants table	Server administration
VERSION_TOKEN_ADMIN	See global_grants table	Server administration
XA_RECOVER_ADMIN	See global_grants table	Server administration

It is a good idea to grant to an account only those privileges that it needs. You should exercise particular caution in granting the FILE and administrative privileges:

- The FILE privilege can be abused to read into a database table any files that the MySQL server can read on the server host. This includes all world-readable files and files in the server's data directory. The table can then be accessed using SELECT to transfer its contents to the client host.
- The GRANT OPTION privilege enables users to give their privileges to other users. Two users that have different privileges and with the GRANT OPTION privilege are able to combine privileges.
- The ALTER privilege may be used to subvert the privilege system by renaming tables.
- The SHUTDOWN privilege can be abused to deny service to other users entirely by terminating the server.
- The PROCESS privilege can be used to view the plain text of currently executing statements, including statements that set or change passwords.
- The SUPER privilege can be used to terminate other sessions or change how the server operates.
- Privileges granted for the mysql database itself can be used to change passwords and other access privilege information. Passwords are stored encrypted, so a malicious user cannot simply read them to know the plain text password. However, a user with write access to the user table authentication_string column can change an account's password, and then connect to the MySQL server using that account.

  The SELECT privilege is also needed for tables or views being used with EXPLAIN, including any underlying tables of views.

The following sections provide general descriptions of the static and dynamic privileges available in MySQL. (For information about the differences between these two types of privileges, see Section 6.2.2, “Static Versus Dynamic Privileges”.) Particular SQL statements might have more specific privilege requirements than indicated here. If so, the description for the statement in question provides the details.

- Static Privileges
- Dynamic Privileges

Static Privileges

Static privileges are built in to the server, in contrast to dynamic privileges, which are defined at runtime. The following list describes the static privileges available in MySQL.

- The ALL or ALL PRIVILEGES privilege specifier is shorthand. It stands for “all privileges available at a given privilege level” (except GRANT OPTION). For example, granting ALL at the global or table level grants all global privileges or all table-level privileges.
- The ALTER privilege enables use of the ALTER TABLE statement to change the structure of tables. ALTER TABLE also requires the CREATE and INSERT privileges. Renaming a table requires ALTER and DROP on the old table, CREATE, and INSERT on the new table.

- The `ALTER ROUTINE` privilege is needed to alter or drop stored routines (procedures and functions).
- The `CREATE` privilege enables creation of new databases and tables.
- The `CREATE ROLE` privilege enables use of the `CREATE ROLE` statement. (The `CREATE USER` privilege also enables use of the `CREATE ROLE` statement.)
- The `CREATE ROUTINE` privilege is needed to create stored routines (procedures and functions).
- The `CREATE TABLESPACE` privilege is needed to create, alter, or drop tablespaces and log file groups.
- The `CREATE TEMPORARY TABLES` privilege enables the creation of temporary tables using the `CREATE TEMPORARY TABLE` statement.

  After a session has created a temporary table, the server performs no further privilege checks on the table. The creating session can perform any operation on the table, such as `DROP TABLE`, `INSERT`, `UPDATE`, or `SELECT`. For more information, see Section 13.1.15.3, “CREATE TEMPORARY TABLE Syntax”.
- The `CREATE USER` privilege enables use of the `ALTER USER`, `CREATE ROLE`, `CREATE USER`, `DROP ROLE`, `DROP USER`, `RENAME USER`, and `REVOKE ALL PRIVILEGES` statements.
- The `CREATE VIEW` privilege enables use of the `CREATE VIEW` statement.
- The `DELETE` privilege enables rows to be deleted from tables in a database.
- The `DROP` privilege enables you to drop (remove) existing databases, tables, and views. The `DROP` privilege is required in order to use the statement `ALTER TABLE ... DROP PARTITION` on a partitioned table. The `DROP` privilege is also required for `TRUNCATE TABLE`. *If you grant the `DROP` privilege for the `mysql` database to a user, that user can drop the database in which the MySQL access privileges are stored.*
- The `DROP ROLE` privilege enables use of the `DROP ROLE` statement. (The `CREATE USER` privilege also enables use of the `DROP ROLE` statement.)
- The `EVENT` privilege is required to create, alter, drop, or see events for the Event Scheduler.
- The `EXECUTE` privilege is required to execute stored routines (procedures and functions).
- The `FILE` privilege gives you permission to read and write files on the server host using the `LOAD DATA INFILE` and `SELECT ... INTO OUTFILE` statements and the `LOAD_FILE()` function. A user who has the `FILE` privilege can read any file on the server host that is either world-readable or readable by the MySQL server. (This implies the user can read any file in any database directory, because the server can access any of those files.) The `FILE` privilege also enables the user to create new files in any directory where the MySQL server has write access. This includes the server's data directory containing the files that implement the privilege tables. As a security measure, the server will not overwrite existing files. The `FILE` privilege is required to use the `DATA DIRECTORY` or `INDEX DIRECTORY` table option for the `CREATE TABLE` statement.

  To limit the location in which files can be read and written, set the `secure_file_priv` system to a specific directory. See Section 5.1.5, “Server System Variables”.
- The `GRANT OPTION` privilege enables you to give to other users or remove from other users those privileges that you yourself possess.
- The `INDEX` privilege enables you to create or drop (remove) indexes. `INDEX` applies to existing tables. If you have the `CREATE` privilege for a table, you can include index definitions in the `CREATE TABLE` statement.
- The `INSERT` privilege enables rows to be inserted into tables in a database. `INSERT` is also required for the `ANALYZE TABLE`, `OPTIMIZE TABLE`, and `REPAIR TABLE` table-maintenance statements.

- The `LOCK TABLES` privilege enables the use of explicit `LOCK TABLES` statements to lock tables for which you have the `SELECT` privilege. This includes the use of write locks, which prevents other sessions from reading the locked table.

- The `PROCESS` privilege pertains to display of information about the threads executing within the server (that is, information about the statements being executed by sessions). The privilege enables use of `SHOW PROCESSLIST` or `mysqladmin processlist` to see threads belonging to other accounts; you can always see your own threads. The `PROCESS` privilege also enables use of `SHOW ENGINE`.

- The `PROXY` privilege enables a user to impersonate or become known as another user. See Section 6.3.11, “Proxy Users”.

- The creation of a foreign key constraint requires the `REFERENCES` privilege for the parent table.

- The `RELOAD` privilege enables use of the `FLUSH` statement. It also enables `mysqladmin` commands that are equivalent to `FLUSH` operations: `flush-hosts`, `flush-logs`, `flush-privileges`, `flush-status`, `flush-tables`, `flush-threads`, `refresh`, and `reload`.

  The `reload` command tells the server to reload the grant tables into memory. `flush-privileges` is a synonym for `reload`. The `refresh` command closes and reopens the log files and flushes all tables. The other `flush-`*`xxx`* commands perform functions similar to `refresh`, but are more specific and may be preferable in some instances. For example, if you want to flush just the log files, `flush-logs` is a better choice than `refresh`.

- The `REPLICATION CLIENT` privilege enables the use of the `SHOW MASTER STATUS`, `SHOW SLAVE STATUS`, and `SHOW BINARY LOGS` statements.

- The `REPLICATION SLAVE` privilege should be granted to accounts that are used by slave servers to connect to the current server as their master. Without this privilege, the slave cannot request updates that have been made to databases on the master server.

- The `SELECT` privilege enables you to select rows from tables in a database. `SELECT` statements require the `SELECT` privilege only if they actually retrieve rows from a table. Some `SELECT` statements do not access tables and can be executed without permission for any database. For example, you can use `SELECT` as a simple calculator to evaluate expressions that make no reference to tables:

  ```
  SELECT 1+1;
  SELECT PI()*2;
  ```

  The `SELECT` privilege is also needed for other statements that read column values. For example, `SELECT` is needed for columns referenced on the right hand side of *`col_name`*=*`expr`* assignment in `UPDATE` statements or for columns named in the `WHERE` clause of `DELETE` or `UPDATE` statements.

- The `SHOW DATABASES` privilege enables the account to see database names by issuing the `SHOW DATABASE` statement. Accounts that do not have this privilege see only databases for which they have some privileges, and cannot use the statement at all if the server was started with the `--skip-show-database` option. Note that *any* global privilege is a privilege for the database.

- The `SHOW VIEW` privilege enables use of the `SHOW CREATE VIEW` statement. This privilege is also needed for views being used with `EXPLAIN`.

- The `SHUTDOWN` privilege enables use of the `SHUTDOWN` statement, the `mysqladmin shutdown` command, and the `mysql_shutdown()` C API function.

- The `SUPER` privilege enables these operations and server behaviors:

  - Enables configuration changes by modifying or persisting global system variables. For some system variables, setting the session value also requires the `SUPER` privilege; if so, it is indicated in the variable description. Examples include `binlog_format`, `sql_log_bin`, and `sql_log_off`.

The corresponding dynamic privilege is `SYSTEM_VARIABLES_ADMIN`.

- Enables starting and stopping replication on slave servers, including Group Replication.

  The corresponding dynamic privilege is `REPLICATION_SLAVE_ADMIN` for regular replication, `GROUP_REPLICATION_ADMIN` for Group Replication.

- Enables use of the `CHANGE MASTER TO` and `CHANGE REPLICATION FILTER` statements.

  The corresponding dynamic privilege is `REPLICATION_SLAVE_ADMIN`.

- Enables binary log control by means of the `PURGE BINARY LOGS` and `BINLOG` statements.

  The corresponding dynamic privilege is `BINLOG_ADMIN`.

- Enables setting the effective authorization ID when executing a view or stored program. A user with this privilege can specify any account in the `DEFINER` attribute of a view or stored program.

  The corresponding dynamic privilege is `SET_USER_ID`.

- Enables use of the `CREATE SERVER`, `ALTER SERVER`, and `DROP SERVER` statements.

- Enables use of the `mysqladmin debug` command.

- Enables `InnoDB` key rotation.

  The corresponding dynamic privilege is `ENCRYPTION_KEY_ADMIN`.

- Enables reading the DES key file by the `DES_ENCRYPT()` function.

- Enables execution of Version Tokens user-defined functions.

  The corresponding dynamic privilege is `VERSION_TOKEN_ADMIN`.

- Enables nonempty `<graphml>` element content in the result from the `ROLES_GRAPHML()` function.

  The corresponding dynamic privilege is `ROLE_ADMIN`.

- Enables control over client connections not permitted to non-`SUPER` accounts:

  - Enables use of the `KILL` statement or `mysqladmin kill` command to kill threads belonging to other accounts. (You can always kill your own threads.)

  - The server accepts one connection from a `SUPER` client even if the connection limit controlled by the `max_connections` system variable is reached.

  - Updates can be performed even when the `read_only` system variable is enabled. This applies to table updates and use of account-management statements such as `GRANT` and `REVOKE`.

    The corresponding dynamic privilege is `CONNECTION_ADMIN`.

  - The server does not execute `init_connect` system variable content when `SUPER` clients connect.

    The corresponding dynamic privilege is `CONNECTION_ADMIN`.

  - A server in offline mode (`offline_mode` enabled) does not terminate `SUPER` client connections at the next client request, and accepts new connections from `SUPER` clients.

    The corresponding dynamic privilege is `CONNECTION_ADMIN`.

You may also need the SUPER privilege to create or alter stored functions if binary logging is enabled, as described in Section 23.7, “Binary Logging of Stored Programs”.

- The TRIGGER privilege enables trigger operations. You must have this privilege for a table to create, drop, execute, or display triggers for that table.

  When a trigger is activated (by a user who has privileges to execute INSERT, UPDATE, or DELETE statements for the table associated with the trigger), trigger execution requires that the user who defined the trigger still have the TRIGGER privilege.

- The UPDATE privilege enables rows to be updated in tables in a database.

- The USAGE privilege specifier stands for “no privileges.” It is used at the global level with GRANT to modify account attributes such as resource limits or SSL characteristics without naming specific account privileges. SHOW GRANTS displays USAGE to indicate that an account has no privileges at a privilege level.

Dynamic Privileges

Dynamic privileges are defined at runtime, in contrast to static privileges, which are built in to the server. The following list describes the dynamic privileges available in MySQL.

- AUDIT_ADMIN: Enables audit log configuration.

- BINLOG_ADMIN: Enables binary log control by means of the PURGE BINARY LOGS and BINLOG statements. Defined at server startup.

- CONNECTION_ADMIN: Enables setting system variables related to client connections, or circumventing restrictions related to client connections. Defined at server startup.

  CONNECTION_ADMIN applies to the effects of these system variables:

  - init_connect: The server does not execute init_connect system variable content when CONNECTION_ADMIN clients connect.

  - offline_mode: A server in offline mode (offline_mode enabled) does not terminate CONNECTION_ADMIN client connections at the next client request, and accepts new connections from CONNECTION_ADMIN clients.

  - read_only: Updates can be performed even when the read_only system variable is enabled. This applies to table updates and use of account-management statements such as GRANT and REVOKE.

- ENCRYPTION_KEY_ADMIN: Enables InnoDB encryption key rotation. Defined at server startup.

- FIREWALL_ADMIN: Enables a user to administer firewall rules for any user.

- FIREWALL_USER: Enables users to update their own firewall rules.

- GROUP_REPLICATION_ADMIN: On a slave server, enables starting and stopping Group Replication. Defined at server startup.

- REPLICATION_SLAVE_ADMIN: On a slave server, enables connecting to and disconnecting from the master server, starting and stopping replication, and use of the CHANGE MASTER TO and CHANGE REPLICATION FILTER statements. Defined at server startup. This privilege does not apply to Group Replication; use GROUP_REPLICATION_ADMIN for that.

- ROLE_ADMIN: Enables use of the WITH ADMIN OPTION clause of the GRANT statement. Enables nonempty <graphml> element content in the result from the ROLES_GRAPHML() function. Defined at server startup.

- `SET_USER_ID`: Enables setting the effective authorization ID when executing a view or stored program. A user with this privilege can specify any account in the `DEFINER` attribute of a view or stored program. Defined at server startup.
- `SYSTEM_VARIABLES_ADMIN`: Enables configuration changes by modifying or persisting global system variables. For some system variables, setting the session value also requires the `SYSTEM_VARIABLES_ADMIN` privilege; if so, it is indicated in the variable description. Examples include `binlog_format`, `sql_log_bin`, and `sql_log_off`. Defined at server startup.
- `VERSION_TOKEN_ADMIN`: Enables execution of Version Tokens user-defined functions. Defined by the `version_tokens` plugin; see Section 5.6.5, “Version Tokens”.
- `XA_RECOVER_ADMIN`: Enables execution of the `XA RECOVER` statement; see Section 13.3.7.1, “XA Transaction SQL Syntax”. Defined at server startup.

  This privilege was added in MySQL 8.0.2. Previously, any user could execute the `XA RECOVER` statement to discover the XID values for outstanding prepared XA transactions, possibly leading to commit or rollback of an XA transaction by a user other than the one who started it. As of 8.0.2, `XA RECOVER` is permitted only to users who have the `XA_RECOVER_ADMIN` privilege, which is expected to be granted only to administrative users who have need for it. This might be the case, for example, for administrators of an XA application if it has crashed and it is necessary to find outstanding transactions started by the application so they can be rolled back. This privilege requirement prevents users from discovering the XID values for outstanding prepared XA transactions other than their own. It does not affect normal commit or rollback of an XA transaction because the user who started it knows its XID.

6.2.2 Static Versus Dynamic Privileges

MySQL supports static and dynamic privileges:

- Static privileges are built in to the server. They are always available to be granted to user accounts and cannot be unregistered.
- Dynamic privileges can be registered and unregistered at runtime. This affects their availability: A dynamic privilege that has not been registered cannot be granted.

For example, the `SELECT` and `INSERT` privileges are static and always available, whereas a dynamic privilege becomes available only if the server component that implements it has been enabled.

The remainder of this section describes how dynamic privileges work in MySQL. The discussion uses the term “components” but applies equally to plugins.

Note

Server administrators should be aware of which server components define dynamic privileges. For MySQL distributions, documentation of components that define dynamic privileges describes those privileges.

Third-party components may also define dynamic privileges; an administrator should understand those privileges and not install components that might conflict or compromise server operation. For example, one component conflicts with another if both define a privilege with the same name. Component developers can reduce the likelihood of this occurrence by choosing privilege names having a prefix based on the component name.

The server maintains the set of registered dynamic privileges internally in memory. Unregistration occurs at server shutdown.

Normally, a server component that defines dynamic privileges registers them when it is installed, during its initialization sequence. When uninstalled, a server component does not unregister its registered

dynamic privileges. (This is current practice, not a requirement. That is, components could, but do not, unregister at any time privileges they register.)

No warning or error occurs for attempts to register an already registered dynamic privilege. Consider the following sequence of statements:

```
INSTALL COMPONENT 'my_component';
UNINSTALL COMPONENT 'my_component';
INSTALL COMPONENT 'my_component';
```

The first `INSTALL COMPONENT` statement registers any privileges defined by server component `my_component`, but `UNINSTALL COMPONENT` does not unregister them. For the second `INSTALL COMPONENT` statement, the component privileges it registers are found to be already registered, but no warnings or errors occur.

Dynamic privileges apply only at the global level. The server stores information about current assignments of dynamic privileges to user accounts in the `mysql.global_grants` system table:

- The server automatically registers privileges named in `global_grants` during server startup (unless the `--skip-grant-tables` option is given).
- The `GRANT` and `REVOKE` statements modify the contents of `global_grants`.
- Dynamic privilege assignments listed in `global_grants` are persistent. They are not removed at server shutdown.

Example: The following statement grants to user `u1` the privileges required to control replication (including Group Replication) on a slave server, and to modify system variables:

```
GRANT REPLICATION_SLAVE_ADMIN, GROUP_REPLICATION_ADMIN, BINLOG_ADMIN
ON *.* TO 'u1'@'localhost';
```

Granted dynamic privileges appear in the output from the `SHOW GRANTS` statement and the `INFORMATION_SCHEMA USER_PRIVILEGES` table.

For `GRANT` and `REVOKE` at the global level, any named privileges not recognized as static are checked against the current set of registered dynamic privileges and granted if found. Otherwise, an error occurs to indicate an unknown privilege identifier.

For `GRANT` and `REVOKE` the meaning of `ALL [PRIVILEGES]` at the global level includes all static global privileges, as well as all currently registered dynamic privileges:

- `GRANT ALL` at the global level grants all static global privileges and all currently registered dynamic privileges. A dynamic privilege registered subsequent to execution of the `GRANT` statement is not granted retroactively to any account.
- `REVOKE ALL` at the global level revokes all granted static global privileges and all granted dynamic privileges.

The `FLUSH PRIVILEGES` statement reads the `global_grants` table for dynamic privilege assignments and registers any unregistered privileges found there.

For descriptions of the dynamic privileges provided by MySQL Server and server components included in MySQL distributions, see Section 6.2.1, “Privileges Provided by MySQL”.

Migrating Accounts from SUPER to Dynamic Privileges

In MySQL 8.0, many operations that previously required the `SUPER` privilege are also associated with a dynamic privilege of more limited scope. (For descriptions of these privileges, see Section 6.2.1,

"Privileges Provided by MySQL".) Each such operation can be permitted to an account by granting the associated dynamic privilege rather than SUPER. This change improves security by enabling DBAs to avoid granting SUPER and tailor user privileges more closely to the operations permitted. SUPER is now deprecated and will be removed in a future version of MySQL.

When removal of SUPER occurs, operations that formerly required SUPER will fail unless accounts granted SUPER are migrated to the appropriate dynamic privileges. Use the following instructions to accomplish that goal so that accounts are ready prior to SUPER removal:

1. Execute this query to identify accounts that are granted SUPER:

   ```
   SELECT GRANTEE FROM INFORMATION_SCHEMA.USER_PRIVILEGES
   WHERE PRIVILEGE_TYPE = 'SUPER';
   ```

2. For each account identified by the preceding query, determine the operations for which it needs SUPER. Then grant the dynamic privileges corresponding to those operations, and revoke SUPER.

   For example, if 'u1'@'localhost' requires SUPER for binary log purging and system variable modification, these statements make the required changes to the account:

   ```
   GRANT BINLOG_ADMIN, SYSTEM_VARIABLES_ADMIN ON *.* TO 'u1'@'localhost';
   REVOKE SUPER ON *.* FROM 'u1'@'localhost';
   ```

   After you have modified all applicable accounts, the INFORMATION_SCHEMA query in the first step should produce an empty result set.

6.2.3 Grant Tables

The mysql system database includes several grant tables that contain information about user accounts and the privileges held by them. This section describes those tables. For information about other tables in the system database, see Section 5.3, "The mysql System Database".

Normally, to manipulate the contents of grant tables, you modify them indirectly by using account-management statements such as CREATE USER, GRANT, and REVOKE to set up accounts and control the privileges available to each one. See Section 13.7.1, "Account Management Statements". The discussion here describes the underlying structure of the grant tables and how the server uses their contents when interacting with clients.

Note

Direct modification of grant tables using statements such as INSERT, UPDATE, or DELETE is discouraged and done at your own risk. The server is free to ignore rows that become malformed as a result of such modifications.

For any operation that modifies a grant table, the server checks whether the table has the expected structure and produces an error if not. mysql_upgrade must be run to update the tables to the expected structure.

These mysql database tables contain grant information:

- user: User accounts, global privileges, and other non-privilege columns
- global_grants: Assignments of dynamic global privileges to users; see Section 6.2.2, "Static Versus Dynamic Privileges".
- db: Database-level privileges
- tables_priv: Table-level privileges
- columns_priv: Column-level privileges

- `procs_priv`: Stored procedure and function privileges
- `proxies_priv`: Proxy-user privileges
- `default_roles`: Default user roles
- `role_edges`: Edges for role subgraphs

In MySQL 8.0, grant tables use the `InnoDB` storage engine and are transactional. Before MySQL 8.0, grant tables used the `MyISAM` storage engine and were nontransactional. This change of grant table storage engine enables an accompanying change to the behavior of account-management statements such as `CREATE USER` or `GRANT`. Previously, an account-management statement that named multiple users could succeed for some users and fail for others. Now, each statement is transactional and either succeeds for all named users or rolls back and has no effect if any error occurs.

Each grant table contains scope columns and privilege columns:

- Scope columns determine the scope of each row in the tables; that is, the context in which the row applies. For example, a `user` table row with `Host` and `User` values of `'thomas.loc.gov'` and `'bob'` applies to authenticating connections made to the server from the host `thomas.loc.gov` by a client that specifies a user name of `bob`. Similarly, a `db` table row with `Host`, `User`, and `Db` column values of `'thomas.loc.gov'`, `'bob'` and `'reports'` applies when `bob` connects from the host `thomas.loc.gov` to access the `reports` database. The `tables_priv` and `columns_priv` tables contain scope columns indicating tables or table/column combinations to which each row applies. The `procs_priv` scope columns indicate the stored routine to which each row applies.
- Privilege columns indicate which privileges a table row grants; that is, which operations it permits to be performed. The server combines the information in the various grant tables to form a complete description of a user's privileges. Section 6.2.7, “Access Control, Stage 2: Request Verification”, describes the rules for this.

The server uses the grant tables in the following manner:

- The `user` table scope columns determine whether to reject or permit incoming connections. For permitted connections, any privileges granted in the `user` table indicate the user's global privileges. Any privileges granted in this table apply to *all* databases on the server.

  Caution

  Because any global privilege is considered a privilege for all databases, any global privilege enables a user to see all database names with `SHOW DATABASES` or by examining the `SCHEMATA` table of `INFORMATION_SCHEMA`.

- The `db` table scope columns determine which users can access which databases from which hosts. The privilege columns determine the permitted operations. A privilege granted at the database level applies to the database and to all objects in the database, such as tables and stored programs.
- The `tables_priv` and `columns_priv` tables are similar to the `db` table, but are more fine-grained: They apply at the table and column levels rather than at the database level. A privilege granted at the table level applies to the table and to all its columns. A privilege granted at the column level applies only to a specific column.
- The `procs_priv` table applies to stored routines (procedures and functions). A privilege granted at the routine level applies only to a single procedure or function.
- The `proxies_priv` table indicates which users can act as proxies for other users and whether a user can grant the `PROXY` privilege to other users.

The server uses the `user` and `db` tables in the `mysql` database at both the first and second stages of access control (see Section 6.2, “The MySQL Access Privilege System”). The columns in the `user` and `db` tables are shown here.

Table 6.4 user and db Table Columns

Table Name	user	db
Scope columns	Host	Host
	User	Db
Privilege columns	Select_priv	Select_priv
	Insert_priv	Insert_priv
	Update_priv	Update_priv
	Delete_priv	Delete_priv
	Index_priv	Index_priv
	Alter_priv	Alter_priv
	Create_priv	Create_priv
	Drop_priv	Drop_priv
	Grant_priv	Grant_priv
	Create_view_priv	Create_view_priv
	Show_view_priv	Show_view_priv
	Create_routine_priv	Create_routine_priv
	Alter_routine_priv	Alter_routine_priv
	Execute_priv	Execute_priv
	Trigger_priv	Trigger_priv
	Event_priv	Event_priv
	Create_tmp_table_priv	Create_tmp_table_priv
	Lock_tables_priv	Lock_tables_priv
	References_priv	References_priv
	Reload_priv	
	Shutdown_priv	
	Process_priv	
	File_priv	
	Show_db_priv	
	Super_priv	
	Repl_slave_priv	
	Repl_client_priv	
	Create_user_priv	
	Create_tablespace_priv	
	Create_role_priv	
	Drop_role_priv	
Security columns	ssl_type	
	ssl_cipher	
	x509_issuer	
	x509_subject	
	plugin	
	authentication_string	
	password_expired	

Table Name	user	db
	password_last_changed	
	password_lifetime	
	account_locked	
Resource control columns	max_questions	
	max_updates	
	max_connections	
	max_user_connections	

The user table plugin and authentication_string columns store authentication plugin and credential information.

The server uses the plugin named in the plugin column of an account row to authenticate connection attempts for the account.

The plugin column must be nonempty. At startup, and at runtime when FLUSH PRIVILEGES is executed, the server checks user table rows. For any row with an empty plugin column, the server writes a warning to the error log of this form:

```
[Warning] User entry 'user_name'@'host_name' has an empty plugin
value. The user will be ignored and no one can login with this user
anymore.
```

The password_expired column permits DBAs to expire account passwords and require users to reset their password. The default password_expired value is 'N', but can be set to 'Y' with the ALTER USER statement. After an account's password has been expired, all operations performed by the account in subsequent connections to the server result in an error until the user issues an ALTER USER statement to establish a new account password.

It is possible after password expiration to “reset” a password by setting it to its current value. As a matter of good policy, it is preferable to choose a different password.

password_last_changed is a TIMESTAMP column indicating when the password was last changed. The value is non-NULL only for accounts that use MySQL built-in authentication methods (accounts that use an authentication plugin of mysql_native_password or sha256_password). The value is NULL for other accounts, such as those authenticated using an external authentication system.

password_last_changed is updated by the CREATE USER, ALTER USER, and SET PASSWORD statements, and by GRANT statements that create an account or change an account password.

password_lifetime indicates the account password lifetime, in days. If the password is past its lifetime (assessed using the password_last_changed column), the server considers the password expired when clients connect using the account. A value of *N* greater than zero means that the password must be changed every *N* days. A value of 0 disables automatic password expiration. If the value is NULL (the default), the global expiration policy applies, as defined by the default_password_lifetime system variable.

account_locked indicates whether the account is locked (see Section 6.3.12, “User Account Locking”).

During the second stage of access control, the server performs request verification to ensure that each client has sufficient privileges for each request that it issues. In addition to the user and db grant tables, the server may also consult the tables_priv and columns_priv tables for requests that involve tables. The latter tables provide finer privilege control at the table and column levels. They have the columns shown in the following table.

Table 6.5 tables_priv and columns_priv Table Columns

Table Name	`tables_priv`	`columns_priv`
Scope columns	`Host`	`Host`
	`Db`	`Db`
	`User`	`User`
	`Table_name`	`Table_name`
		`Column_name`
Privilege columns	`Table_priv`	`Column_priv`
	`Column_priv`	
Other columns	`Timestamp`	`Timestamp`
	`Grantor`	

The `Timestamp` and `Grantor` columns are set to the current timestamp and the `CURRENT_USER` value, respectively, but are otherwise unused.

For verification of requests that involve stored routines, the server may consult the `procs_priv` table, which has the columns shown in the following table.

Table 6.6 procs_priv Table Columns

Table Name	`procs_priv`
Scope columns	`Host`
	`Db`
	`User`
	`Routine_name`
	`Routine_type`
Privilege columns	`Proc_priv`
Other columns	`Timestamp`
	`Grantor`

The `Routine_type` column is an `ENUM` column with values of `'FUNCTION'` or `'PROCEDURE'` to indicate the type of routine the row refers to. This column enables privileges to be granted separately for a function and a procedure with the same name.

The `Timestamp` and `Grantor` columns are unused.

The `proxies_priv` table records information about proxy accounts. It has these columns:

- `Host`, `User`: The proxy account; that is, the account that has the `PROXY` privilege for the proxied account.
- `Proxied_host`, `Proxied_user`: The proxied account.
- `Grantor`, `Timestamp`: Unused.
- `With_grant`: Whether the proxy account can grant the `PROXY` privilege to other accounts.

For an account to be able to grant the `PROXY` privilege to other accounts, it must have a row in the `proxies_priv` table with `With_grant` set to 1 and `Proxied_host` and `Proxied_user` set to indicate the account or accounts for which the privilege can be granted. For example, the `'root'@'localhost'` account created during MySQL installation has a row in the `proxies_priv` table that enables granting the `PROXY` privilege for `''@''`, that is, for all users and all hosts. This enables `root` to set up proxy users, as well as to delegate to other accounts the authority to set up proxy users. See Section 6.3.11, “Proxy Users”.

The `global_grants` table lists current assignments of dynamic privileges to user accounts. These privileges are global. The table has these columns:

- `USER`, `HOST`: The user name and host name of the account to which the privilege is granted.
- `PRIV`: The privilege name.
- `WITH_GRANT_OPTION`: Whether the account can grant the privilege to other accounts.

The `default_roles` table lists default user roles. It has these columns:

- `HOST`, `USER`: The account or role to which the default role applies.
- `DEFAULT_ROLE_HOST`, `DEFAULT_ROLE_USER`: The default role.

The `role_edges` table lists edges for role subgraphs. It has these columns:

- `FROM_HOST`, `FROM_USER`: The account that is granted a role.
- `TO_HOST`, `TO_USER`: The role that is granted to the account.
- `WITH_ADMIN_OPTION`: Whether the account can grant the role to and revoke it from other accounts by using `WITH ADMIN OPTION`.

Scope columns in the grant tables contain strings. The default value for each is the empty string. The following table shows the number of characters permitted in each column.

Table 6.7 Grant Table Scope Column Lengths

Column Name	Maximum Permitted Characters
`Host`, `Proxied_host`	60
`User`, `Proxied_user`	32
`Db`	64
`Table_name`	64
`Column_name`	64
`Routine_name`	64

For access-checking purposes, comparisons of `User`, `Proxied_user`, `authentication_string`, `Db`, and `Table_name` values are case sensitive. Comparisons of `Host`, `Proxied_host`, `Column_name`, and `Routine_name` values are not case sensitive.

The `user` and `db` tables list each privilege in a separate column that is declared as `ENUM('N','Y') DEFAULT 'N'`. In other words, each privilege can be disabled or enabled, with the default being disabled.

The `tables_priv`, `columns_priv`, and `procs_priv` tables declare the privilege columns as `SET` columns. Values in these columns can contain any combination of the privileges controlled by the table. Only those privileges listed in the column value are enabled.

Table 6.8 Set-Type Privilege Column Values

Table Name	Column Name	Possible Set Elements
`tables_priv`	`Table_priv`	`'Select', 'Insert', 'Update', 'Delete', 'Create', 'Drop', 'Grant', 'References', 'Index', 'Alter', 'Create View', 'Show view', 'Trigger'`
`tables_priv`	`Column_priv`	`'Select', 'Insert', 'Update', 'References'`

Table Name	Column Name	Possible Set Elements
columns_priv	Column_priv	'Select', 'Insert', 'Update', 'References'
procs_priv	Proc_priv	'Execute', 'Alter Routine', 'Grant'

Only the user table specifies administrative privileges, such as RELOAD and SHUTDOWN. Administrative operations are operations on the server itself and are not database-specific, so there is no reason to list these privileges in the other grant tables. Consequently, the server need consult only the user table to determine whether a user can perform an administrative operation.

The FILE privilege also is specified only in the user table. It is not an administrative privilege as such, but a user's ability to read or write files on the server host is independent of the database being accessed.

The server reads the contents of the grant tables into memory when it starts. You can tell it to reload the tables by issuing a FLUSH PRIVILEGES statement or executing a mysqladmin flush-privileges or mysqladmin reload command. Changes to the grant tables take effect as indicated in Section 6.2.8, “When Privilege Changes Take Effect”.

When you modify an account, it is a good idea to verify that your changes have the intended effect. To check the privileges for a given account, use the SHOW GRANTS statement. For example, to determine the privileges that are granted to an account with user name and host name values of bob and pc84.example.com, use this statement:

```
SHOW GRANTS FOR 'bob'@'pc84.example.com';
```

To display nonprivilege properties of an account, use SHOW CREATE USER:

```
SHOW CREATE USER 'bob'@'pc84.example.com';
```

6.2.4 Specifying Account Names

MySQL account names consist of a user name and a host name. This enables creation of accounts for users with the same name who can connect from different hosts. This section describes how to write account names, including special values and wildcard rules.

MySQL role names are similar to account names, with some differences described at Section 6.2.5, “Specifying Role Names”.

In SQL statements such as CREATE USER, GRANT, and SET PASSWORD, account names follow these rules:

- Account name syntax is *'user_name'@'host_name'*.
- An account name consisting only of a user name is equivalent to *'user_name'*'@'%'. For example, 'me' is equivalent to 'me'@'%'.
- The user name and host name need not be quoted if they are legal as unquoted identifiers. Quotes are necessary to specify a *user_name* string containing special characters (such as space or -), or a *host_name* string containing special characters or wildcard characters (such as . or %); for example, 'test-user'@'%.com'.
- Quote user names and host names as identifiers or as strings, using either backticks (`), single quotation marks ('), or double quotation marks ("). For string-quoting and identifier-quoting guidelines, see Section 9.1.1, “String Literals”, and Section 9.2, “Schema Object Names”.
- The user name and host name parts, if quoted, must be quoted separately. That is, write 'me'@'localhost', not 'me@localhost'; the latter is actually equivalent to 'me@localhost'@'%'.

- A reference to the `CURRENT_USER` or `CURRENT_USER()` function is equivalent to specifying the current client's user name and host name literally.

MySQL stores account names in grant tables in the `mysql` system database using separate columns for the user name and host name parts:

- The `user` table contains one row for each account. The `User` and `Host` columns store the user name and host name. This table also indicates which global privileges the account has.
- Other grant tables indicate privileges an account has for databases and objects within databases. These tables have `User` and `Host` columns to store the account name. Each row in these tables associates with the account in the `user` table that has the same `User` and `Host` values.
- For access-checking purposes, comparisons of User values are case sensitive. Comparisons of Host values are not case sensitive.

For additional detail about grant table structure, see Section 6.2.3, “Grant Tables”.

User names and host names have certain special values or wildcard conventions, as described following.

The user name part of an account name is either a nonblank value that literally matches the user name for incoming connection attempts, or a blank value (empty string) that matches any user name. An account with a blank user name is an anonymous user. To specify an anonymous user in SQL statements, use a quoted empty user name part, such as `''@'localhost'`.

The host name part of an account name can take many forms, and wildcards are permitted:

- A host value can be a host name or an IP address (IPv4 or IPv6). The name `'localhost'` indicates the local host. The IP address `'127.0.0.1'` indicates the IPv4 loopback interface. The IP address `'::1'` indicates the IPv6 loopback interface.
- The `%` and `_` wildcard characters are permitted in host name or IP address values. These have the same meaning as for pattern-matching operations performed with the `LIKE` operator. For example, a host value of `'%'` matches any host name, whereas a value of `'%.mysql.com'` matches any host in the `mysql.com` domain. `'192.168.1.%'` matches any host in the 192.168.1 class C network.

  Because IP wildcard values are permitted in host values (for example, `'192.168.1.%'` to match every host on a subnet), someone could try to exploit this capability by naming a host `192.168.1.somewhere.com`. To foil such attempts, MySQL does not perform matching on host names that start with digits and a dot. For example, if a host is named `1.2.example.com`, its name never matches the host part of account names. An IP wildcard value can match only IP addresses, not host names.
- For a host value specified as an IPv4 address, a netmask can be given to indicate how many address bits to use for the network number. Netmask notation cannot be used for IPv6 addresses.

  The syntax is *`host_ip/netmask`*. For example:

```
CREATE USER 'david'@'192.58.197.0/255.255.255.0';
```

  This enables `david` to connect from any client host having an IP address *`client_ip`* for which the following condition is true:

```
client_ip & netmask = host_ip
```

  That is, for the `CREATE USER` statement just shown:

```
client_ip & 255.255.255.0 = 192.58.197.0
```

IP addresses that satisfy this condition range from `192.58.197.0` to `192.58.197.255`.

A netmask typically begins with bits set to 1, followed by bits set to 0. Examples:

- `192.0.0.0/255.0.0.0`: Any host on the 192 class A network
- `192.168.0.0/255.255.0.0`: Any host on the 192.168 class B network
- `192.168.1.0/255.255.255.0`: Any host on the 192.168.1 class C network
- `192.168.1.1`: Only the host with this specific IP address

The server performs matching of host values in account names against the client host using the value returned by the system DNS resolver for the client host name or IP address. Except in the case that the account host value is specified using netmask notation, the server performs this comparison as a string match, even for an account host value given as an IP address. This means that you should specify account host values in the same format used by DNS. Here are examples of problems to watch out for:

- Suppose that a host on the local network has a fully qualified name of `host1.example.com`. If DNS returns name lookups for this host as `host1.example.com`, use that name in account host values. If DNS returns just `host1`, use `host1` instead.
- If DNS returns the IP address for a given host as `192.168.1.2`, that will match an account host value of `192.168.1.2` but not `192.168.01.2`. Similarly, it will match an account host pattern like `192.168.1.%` but not `192.168.01.%`.

To avoid problems like these, it is advisable to check the format in which your DNS returns host names and addresses. Use values in the same format in MySQL account names.

6.2.5 Specifying Role Names

MySQL role names refer to roles, which are named collections of privileges. For role usage examples, see Section 6.3.4, “Using Roles”.

Role names have syntax and semantics similar to account names (Section 6.2.4, “Specifying Account Names”). Role names differ from account names in these respects:

- The user part of role names cannot be blank. Thus, there is no “anonymous role” analogous to the concept of “anonymous user.”
- As for an account name, omitting the host part of a role name results in a host part of `'%'`. But unlike `'%'` in an account name, a host part of `'%'` in a role name has no wildcard properties. For example, for a name `'me'@'%'` used as a role name, the host part (`'%'`) is just a literal value; it has no “any host” matching property.
- Netmask notation in the host part of a role name has no significance.
- An account name is permitted to be `CURRENT_USER()` in several contexts. A role name is not.

It is possible for a row in the `mysql.user` system table to serve as both an account and a role. In this case, any special user or host name matching properties do not apply in contexts for which the name is used as a role name. For example, you cannot execute the following statement with the expectation that it will set the current session roles using all roles that have a user part of `myrole` and any host name:

```
SET ROLE 'myrole'@'%';
```

Instead, the statement sets the active role for the session to the role with exactly the name `'myrole'@'%'`.

For this reason, role names are often specified using only the user name part and letting the host name part implicitly be `'%'`. Specifying a role with a non-`'%'` host part can be useful if you intend to create a name that works both as a role an as a user account that is permitted to connect from the given host.

6.2.6 Access Control, Stage 1: Connection Verification

When you attempt to connect to a MySQL server, the server accepts or rejects the connection based on these conditions:

- Your identity and whether you can verify your identity by supplying the correct password
- Whether your account is locked or unlocked

The server checks credentials first, then account locking state. A failure for either step causes the server to deny access to you completely. Otherwise, the server accepts the connection, and then enters Stage 2 and waits for requests.

Credential checking is performed using the three `user` table scope columns (`Host`, `User`, and `authentication_string`). Locking state is recorded in the `user` table `account_locked` column. The server accepts the connection only if the `Host` and `User` columns in some `user` table row match the client host name and user name, the client supplies the password specified in that row, and the `account_locked` value is `'N'`. The rules for permissible `Host` and `User` values are given in Section 6.2.4, “Specifying Account Names”. Account locking can be changed with the `ALTER USER` statement.

Your identity is based on two pieces of information:

- The client host from which you connect
- Your MySQL user name

If the `User` column value is nonblank, the user name in an incoming connection must match exactly. If the `User` value is blank, it matches any user name. If the `user` table row that matches an incoming connection has a blank user name, the user is considered to be an anonymous user with no name, not a user with the name that the client actually specified. This means that a blank user name is used for all further access checking for the duration of the connection (that is, during Stage 2).

The `authentication_string` column can be blank. This is not a wildcard and does not mean that any password matches. It means that the user must connect without specifying a password. If the server authenticates a client using a plugin, the authentication method that the plugin implements may or may not use the password in the `authentication_string` column. In this case, it is possible that an external password is also used to authenticate to the MySQL server.

Nonblank `authentication_string` values in the `user` table represent encrypted passwords. MySQL does not store passwords in cleartext form for anyone to see. Rather, the password supplied by a user who is attempting to connect is encrypted (using the password hashing method implemented by the account authentication plugin). The encrypted password then is used during the connection process when checking whether the password is correct. This is done without the encrypted password ever traveling over the connection. See Section 6.3.1, “User Names and Passwords”.

From MySQL's point of view, the encrypted password is the *real* password, so you should never give anyone access to it. In particular, *do not give nonadministrative users read access to tables in the `mysql` database*.

The following table shows how various combinations of `User` and `Host` values in the `user` table apply to incoming connections.

`User` Value	`Host` Value	Permissible Connections
`'fred'`	`'thomas.loc.gov'`	`fred`, connecting from `thomas.loc.gov`

`User` Value	`Host` Value	Permissible Connections
`''`	`'thomas.loc.gov'`	Any user, connecting from `thomas.loc.gov`
`'fred'`	`'%'`	`fred`, connecting from any host
`''`	`'%'`	Any user, connecting from any host
`'fred'`	`'%.loc.gov'`	`fred`, connecting from any host in the `loc.gov` domain
`'fred'`	`'x.y.%'`	`fred`, connecting from `x.y.net`, `x.y.com`, `x.y.edu`, and so on; this is probably not useful
`'fred'`	`'192.168.10.177'`	`fred`, connecting from the host with IP address `192.168.10.177`
`'fred'`	`'192.168.10.%'`	`fred`, connecting from any host in the `192.168.10` class C subnet
`'fred'`	`'192.168.10.0/255.255.255.0'`	Same as previous example

It is possible for the client host name and user name of an incoming connection to match more than one row in the `user` table. The preceding set of examples demonstrates this: Several of the entries shown match a connection from `thomas.loc.gov` by `fred`.

When multiple matches are possible, the server must determine which of them to use. It resolves this issue as follows:

- Whenever the server reads the `user` table into memory, it sorts the rows.
- When a client attempts to connect, the server looks through the rows in sorted order.
- The server uses the first row that matches the client host name and user name.

The server uses sorting rules that order rows with the most-specific `Host` values first. Literal host names and IP addresses are the most specific. (The specificity of a literal IP address is not affected by whether it has a netmask, so `192.168.1.13` and `192.168.1.0/255.255.255.0` are considered equally specific.) The pattern `'%'` means “any host” and is least specific. The empty string `''` also means “any host” but sorts after `'%'`. Rows with the same `Host` value are ordered with the most-specific `User` values first (a blank `User` value means “any user” and is least specific). For rows with equally-specific `Host` and `User` values, the order is indeterminate.

To see how this works, suppose that the `user` table looks like this:

```
+-----------+----------+-
| Host      | User     | ...
+-----------+----------+-
| %         | root     | ...
| %         | jeffrey  | ...
| localhost | root     | ...
| localhost |          | ...
+-----------+----------+-
```

When the server reads the table into memory, it sorts the rows using the rules just described. The result after sorting looks like this:

```
+-----------+----------+-
| Host      | User     | ...
+-----------+----------+-
| localhost | root     | ...
| localhost |          | ...
| %         | jeffrey  | ...
| %         | root     | ...
+-----------+----------+-
```

When a client attempts to connect, the server looks through the sorted rows and uses the first match found. For a connection from `localhost` by `jeffrey`, two of the rows from the table match: the one with `Host` and `User` values of `'localhost'` and `''`, and the one with values of `'%'` and `'jeffrey'`. The `'localhost'` row appears first in sorted order, so that is the one the server uses.

Here is another example. Suppose that the `user` table looks like this:

```
+---------------+----------+-
| Host          | User     | ...
+---------------+----------+-
| %             | jeffrey  | ...
| thomas.loc.gov |          | ...
+---------------+----------+-
```

The sorted table looks like this:

```
+---------------+----------+-
| Host          | User     | ...
+---------------+----------+-
| thomas.loc.gov |          | ...
| %             | jeffrey  | ...
+---------------+----------+-
```

A connection by `jeffrey` from `thomas.loc.gov` is matched by the first row, whereas a connection by `jeffrey` from any host is matched by the second.

Note

It is a common misconception to think that, for a given user name, all rows that explicitly name that user are used first when the server attempts to find a match for the connection. This is not true. The preceding example illustrates this, where a connection from `thomas.loc.gov` by `jeffrey` is first matched not by the row containing `'jeffrey'` as the `User` column value, but by the row with no user name. As a result, `jeffrey` is authenticated as an anonymous user, even though he specified a user name when connecting.

If you are able to connect to the server, but your privileges are not what you expect, you probably are being authenticated as some other account. To find out what account the server used to authenticate you, use the `CURRENT_USER()` function. (See Section 12.14, “Information Functions”.) It returns a value in *`user_name@host_name`* format that indicates the `User` and `Host` values from the matching `user` table row. Suppose that `jeffrey` connects and issues the following query:

```
mysql> SELECT CURRENT_USER();
+----------------+
| CURRENT_USER() |
+----------------+
| @localhost     |
+----------------+
```

The result shown here indicates that the matching `user` table row had a blank `User` column value. In other words, the server is treating `jeffrey` as an anonymous user.

Another way to diagnose authentication problems is to print out the `user` table and sort it by hand to see where the first match is being made.

6.2.7 Access Control, Stage 2: Request Verification

After you establish a connection, the server enters Stage 2 of access control. For each request that you issue through that connection, the server determines what operation you want to perform, then checks

whether you have sufficient privileges to do so. This is where the privilege columns in the grant tables come into play. These privileges can come from any of the `user`, `db`, `tables_priv`, `columns_priv`, or `procs_priv` tables. (You may find it helpful to refer to Section 6.2.3, “Grant Tables”, which lists the columns present in each of the grant tables.)

The `user` table grants privileges that are assigned to you on a global basis and that apply no matter what the default database is. For example, if the `user` table grants you the `DELETE` privilege, you can delete rows from any table in any database on the server host! It is wise to grant privileges in the `user` table only to people who need them, such as database administrators. For other users, you should leave all privileges in the `user` table set to `'N'` and grant privileges at more specific levels only. You can grant privileges for particular databases, tables, columns, or routines.

The `db` table grants database-specific privileges. Values in the scope columns of this table can take the following forms:

- A blank `User` value matches the anonymous user. A nonblank value matches literally; there are no wildcards in user names.
- The wildcard characters `%` and `_` can be used in the `Host` and `Db` columns. These have the same meaning as for pattern-matching operations performed with the `LIKE` operator. If you want to use either character literally when granting privileges, you must escape it with a backslash. For example, to include the underscore character (_) as part of a database name, specify it as `\_` in the `GRANT` statement.
- A `'%'` or blank `Host` value means “any host.”
- A `'%'` or blank `Db` value means “any database.”

The server reads the `db` table into memory and sorts it at the same time that it reads the `user` table. The server sorts the `db` table based on the `Host`, `Db`, and `User` scope columns. As with the `user` table, sorting puts the most-specific values first and least-specific values last, and when the server looks for matching rows, it uses the first match that it finds.

The `tables_priv`, `columns_priv`, and `procs_priv` tables grant table-specific, column-specific, and routine-specific privileges. Values in the scope columns of these tables can take the following forms:

- The wildcard characters `%` and `_` can be used in the `Host` column. These have the same meaning as for pattern-matching operations performed with the `LIKE` operator.
- A `'%'` or blank `Host` value means “any host.”
- The `Db`, `Table_name`, `Column_name`, and `Routine_name` columns cannot contain wildcards or be blank.

The server sorts the `tables_priv`, `columns_priv`, and `procs_priv` tables based on the `Host`, `Db`, and `User` columns. This is similar to `db` table sorting, but simpler because only the `Host` column can contain wildcards.

The server uses the sorted tables to verify each request that it receives. For requests that require administrative privileges such as `SHUTDOWN` or `RELOAD`, the server checks only the `user` table row because that is the only table that specifies administrative privileges. The server grants access if the row permits the requested operation and denies access otherwise. For example, if you want to execute `mysqladmin shutdown` but your `user` table row does not grant the `SHUTDOWN` privilege to you, the server denies access without even checking the `db` table. (It contains no `Shutdown_priv` column, so there is no need to do so.)

For database-related requests (`INSERT`, `UPDATE`, and so on), the server first checks the user's global privileges by looking in the `user` table row. If the row permits the requested operation, access is

granted. If the global privileges in the `user` table are insufficient, the server determines the user's database-specific privileges by checking the `db` table:

The server looks in the `db` table for a match on the `Host`, `Db`, and `User` columns. The `Host` and `User` columns are matched to the connecting user's host name and MySQL user name. The `Db` column is matched to the database that the user wants to access. If there is no row for the `Host` and `User`, access is denied.

After determining the database-specific privileges granted by the `db` table rows, the server adds them to the global privileges granted by the `user` table. If the result permits the requested operation, access is granted. Otherwise, the server successively checks the user's table and column privileges in the `tables_priv` and `columns_priv` tables, adds those to the user's privileges, and permits or denies access based on the result. For stored-routine operations, the server uses the `procs_priv` table rather than `tables_priv` and `columns_priv`.

Expressed in boolean terms, the preceding description of how a user's privileges are calculated may be summarized like this:

```
global privileges
OR (database privileges AND host privileges)
OR table privileges
OR column privileges
OR routine privileges
```

It may not be apparent why, if the global `user` row privileges are initially found to be insufficient for the requested operation, the server adds those privileges to the database, table, and column privileges later. The reason is that a request might require more than one type of privilege. For example, if you execute an `INSERT INTO ... SELECT` statement, you need both the `INSERT` and the `SELECT` privileges. Your privileges might be such that the `user` table row grants one privilege and the `db` table row grants the other. In this case, you have the necessary privileges to perform the request, but the server cannot tell that from either table by itself; the privileges granted by the rows in both tables must be combined.

6.2.8 When Privilege Changes Take Effect

When `mysqld` starts, it reads all grant table contents into memory. The in-memory tables become effective for access control at that point.

If you modify the grant tables indirectly using account-management statements such as `GRANT`, `REVOKE`, `SET PASSWORD`, or `RENAME USER`, the server notices these changes and loads the grant tables into memory again immediately.

If you modify the grant tables directly using statements such as `INSERT`, `UPDATE`, or `DELETE`, your changes have no effect on privilege checking until you either restart the server or tell it to reload the tables. If you change the grant tables directly but forget to reload them, your changes have *no effect* until you restart the server. This may leave you wondering why your changes seem to make no difference!

To tell the server to reload the grant tables, perform a flush-privileges operation. This can be done by issuing a `FLUSH PRIVILEGES` statement or by executing a `mysqladmin flush-privileges` or `mysqladmin reload` command.

A grant table reload affects privileges for each existing client connection as follows:

- Table and column privilege changes take effect with the client's next request.
- Database privilege changes take effect the next time the client executes a `USE db_name` statement.

Note

Client applications may cache the database name; thus, this effect may not be visible to them without actually changing to a different database.

- Global privileges and passwords are unaffected for a connected client. These changes take effect only for subsequent connections.

If the server is started with the `--skip-grant-tables` option, it does not read the grant tables or implement any access control. Anyone can connect and do anything, *which is insecure.* To cause a server thus started to read the tables and enable access checking, flush the privileges.

6.2.9 Troubleshooting Problems Connecting to MySQL

If you encounter problems when you try to connect to the MySQL server, the following items describe some courses of action you can take to correct the problem.

- Make sure that the server is running. If it is not, clients cannot connect to it. For example, if an attempt to connect to the server fails with a message such as one of those following, one cause might be that the server is not running:

```
shell> mysql
ERROR 2003: Can't connect to MySQL server on 'host_name' (111)
shell> mysql
ERROR 2002: Can't connect to local MySQL server through socket
'/tmp/mysql.sock' (111)
```

- It might be that the server is running, but you are trying to connect using a TCP/IP port, named pipe, or Unix socket file different from the one on which the server is listening. To correct this when you invoke a client program, specify a `--port` option to indicate the proper port number, or a `--socket` option to indicate the proper named pipe or Unix socket file. To find out where the socket file is, you can use this command:

```
shell> netstat -ln | grep mysql
```

- Make sure that the server has not been configured to ignore network connections or (if you are attempting to connect remotely) that it has not been configured to listen only locally on its network interfaces. If the server was started with `--skip-networking`, it will not accept TCP/IP connections at all. If the server was started with `--bind-address=127.0.0.1`, it will listen for TCP/IP connections only locally on the loopback interface and will not accept remote connections.

- Check to make sure that there is no firewall blocking access to MySQL. Your firewall may be configured on the basis of the application being executed, or the port number used by MySQL for communication (3306 by default). Under Linux or Unix, check your IP tables (or similar) configuration to ensure that the port has not been blocked. Under Windows, applications such as ZoneAlarm or Windows Firewall may need to be configured not to block the MySQL port.

- The grant tables must be properly set up so that the server can use them for access control. For some distribution types (such as binary distributions on Windows, or RPM distributions on Linux), the installation process initializes the MySQL data directory, including the `mysql` database containing the grant tables. For distributions that do not do this, you must initialize the data directory manually. For details, see Section 2.9, “Postinstallation Setup and Testing”.

  To determine whether you need to initialize the grant tables, look for a `mysql` directory under the data directory. (The data directory normally is named `data` or `var` and is located under your MySQL installation directory.) Make sure that you have a file named `user.MYD` in the `mysql` database directory. If not, initialize the data directory. After doing so and starting the server, test the initial privileges by executing this command:

```
shell> mysql -u root
```

  The server should let you connect without error.

- After a fresh installation, you should connect to the server and set up your users and their access permissions:

```
shell> mysql -u root mysql
```

The server should let you connect with no password if you initialized MySQL using `mysqld --initialize-secure` to not create a password for the initial `root` account (see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”). That is a security risk, so setting the password for the `root` account is something you should do while you're setting up your other MySQL accounts. For instructions on setting the initial password, see Section 2.9.4, “Securing the Initial MySQL Account”.

- If you have updated an existing MySQL installation to a newer version, did you run the `mysql_upgrade` script? If not, do so. The structure of the grant tables changes occasionally when new capabilities are added, so after an upgrade you should always make sure that your tables have the current structure. For instructions, see Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

- If a client program receives the following error message when it tries to connect, it means that the server expects passwords in a newer format than the client is capable of generating:

```
shell> mysql
Client does not support authentication protocol requested
by server; consider upgrading MySQL client
```

- Remember that client programs use connection parameters specified in option files or environment variables. If a client program seems to be sending incorrect default connection parameters when you have not specified them on the command line, check any applicable option files and your environment. For example, if you get `Access denied` when you run a client without any options, make sure that you have not specified an old password in any of your option files!

  You can suppress the use of option files by a client program by invoking it with the `--no-defaults` option. For example:

```
shell> mysqladmin --no-defaults -u root version
```

  The option files that clients use are listed in Section 4.2.6, “Using Option Files”. Environment variables are listed in Section 4.9, “MySQL Program Environment Variables”.

- If you get the following error, it means that you are using an incorrect `root` password:

```
shell> mysqladmin -u root -pxxxx ver
Access denied for user 'root'@'localhost' (using password: YES)
```

  If the preceding error occurs even when you have not specified a password, it means that you have an incorrect password listed in some option file. Try the `--no-defaults` option as described in the previous item.

  For information on changing passwords, see Section 6.3.7, “Assigning Account Passwords”.

  If you have lost or forgotten the `root` password, see Section B.5.3.2, “How to Reset the Root Password”.

- If you change a password by using `SET PASSWORD`, `INSERT`, or `UPDATE`, you must encrypt the password using the `PASSWORD()` function. If you do not use `PASSWORD()` for these statements, the password will not work. For example, the following statement assigns a password, but fails to encrypt it, so the user is not able to connect afterward:

```
SET PASSWORD FOR 'abe'@'host_name' = 'eagle';
```

  Instead, set the password like this:

```
SET PASSWORD FOR 'abe'@'host_name' = PASSWORD('eagle');
```

The `PASSWORD()` function is unnecessary when you specify a password using the `CREATE USER` or `GRANT` statements or the `mysqladmin password` command. Each of those automatically uses `PASSWORD()` to encrypt the password. See Section 6.3.7, “Assigning Account Passwords”, and Section 13.7.1.3, “CREATE USER Syntax”.

- `localhost` is a synonym for your local host name, and is also the default host to which clients try to connect if you specify no host explicitly.

  You can use a `--host=127.0.0.1` option to name the server host explicitly. This will make a TCP/IP connection to the local `mysqld` server. You can also use TCP/IP by specifying a `--host` option that uses the actual host name of the local host. In this case, the host name must be specified in a `user` table row on the server host, even though you are running the client program on the same host as the server.

- The `Access denied` error message tells you who you are trying to log in as, the client host from which you are trying to connect, and whether you were using a password. Normally, you should have one row in the `user` table that exactly matches the host name and user name that were given in the error message. For example, if you get an error message that contains `using password: NO`, it means that you tried to log in without a password.

- If you get an `Access denied` error when trying to connect to the database with `mysql -u user_name`, you may have a problem with the `user` table. Check this by executing `mysql -u root mysql` and issuing this SQL statement:

  ```
  SELECT * FROM user;
  ```

  The result should include a row with the `Host` and `User` columns matching your client's host name and your MySQL user name.

- If the following error occurs when you try to connect from a host other than the one on which the MySQL server is running, it means that there is no row in the `user` table with a `Host` value that matches the client host:

  ```
  Host ... is not allowed to connect to this MySQL server
  ```

  You can fix this by setting up an account for the combination of client host name and user name that you are using when trying to connect.

  If you do not know the IP address or host name of the machine from which you are connecting, you should put a row with `'%'` as the `Host` column value in the `user` table. After trying to connect from the client machine, use a `SELECT USER()` query to see how you really did connect. Then change the `'%'` in the `user` table row to the actual host name that shows up in the log. Otherwise, your system is left insecure because it permits connections from any host for the given user name.

  On Linux, another reason that this error might occur is that you are using a binary MySQL version that is compiled with a different version of the `glibc` library than the one you are using. In this case, you should either upgrade your operating system or `glibc`, or download a source distribution of MySQL version and compile it yourself. A source RPM is normally trivial to compile and install, so this is not a big problem.

- If you specify a host name when trying to connect, but get an error message where the host name is not shown or is an IP address, it means that the MySQL server got an error when trying to resolve the IP address of the client host to a name:

  ```
  shell> mysqladmin -u root -pxxxx -h some_hostname ver
  Access denied for user 'root'@'' (using password: YES)
  ```

If you try to connect as `root` and get the following error, it means that you do not have a row in the `user` table with a `User` column value of `'root'` and that `mysqld` cannot resolve the host name for your client:

```
Access denied for user ''@'unknown'
```

These errors indicate a DNS problem. To fix it, execute `mysqladmin flush-hosts` to reset the internal DNS host cache. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.

Some permanent solutions are:

- Determine what is wrong with your DNS server and fix it.
- Specify IP addresses rather than host names in the MySQL grant tables.
- Put an entry for the client machine name in `/etc/hosts` on Unix or `\windows\hosts` on Windows.
- Start `mysqld` with the `--skip-name-resolve` option.
- Start `mysqld` with the `--skip-host-cache` option.
- On Unix, if you are running the server and the client on the same machine, connect to `localhost`. For connections to `localhost`, MySQL programs attempt to connect to the local server by using a Unix socket file, unless there are connection parameters specified to ensure that the client makes a TCP/IP connection. For more information, see Section 4.2.2, “Connecting to the MySQL Server”.
- On Windows, if you are running the server and the client on the same machine and the server supports named pipe connections, connect to the host name `.` (period). Connections to `.` use a named pipe rather than TCP/IP.

- If `mysql -u root` works but `mysql -h your_hostname -u root` results in `Access denied` (where *`your_hostname`* is the actual host name of the local host), you may not have the correct name for your host in the `user` table. A common problem here is that the `Host` value in the `user` table row specifies an unqualified host name, but your system's name resolution routines return a fully qualified domain name (or vice versa). For example, if you have a row with host `'pluto'` in the `user` table, but your DNS tells MySQL that your host name is `'pluto.example.com'`, the row does not work. Try adding a row to the `user` table that contains the IP address of your host as the `Host` column value. (Alternatively, you could add a row to the `user` table with a `Host` value that contains a wildcard; for example, `'pluto.%'`. However, use of `Host` values ending with `%` is *insecure* and is *not* recommended!)

- If `mysql -u user_name` works but `mysql -u user_name some_db` does not, you have not granted access to the given user for the database named *`some_db`*.

- If `mysql -u user_name` works when executed on the server host, but `mysql -h host_name -u user_name` does not work when executed on a remote client host, you have not enabled access to the server for the given user name from the remote host.

- If you cannot figure out why you get `Access denied`, remove from the `user` table all rows that have `Host` values containing wildcards (rows that contain `'%'` or `'_'` characters). A very common error is to insert a new row with `Host='%'` and `User='some_user'`, thinking that this enables you to specify `localhost` to connect from the same machine. The reason that this does not work is that the default privileges include a row with `Host='localhost'` and `User=''`. Because that row has a `Host` value `'localhost'` that is more specific than `'%'`, it is used in preference to the new row when connecting from `localhost`! The correct procedure is to insert a second row with `Host='localhost'` and `User='some_user'`, or to delete the row with `Host='localhost'` and

`User=''`. After deleting the row, remember to issue a `FLUSH PRIVILEGES` statement to reload the grant tables. See also Section 6.2.6, “Access Control, Stage 1: Connection Verification”.

- If you are able to connect to the MySQL server, but get an `Access denied` message whenever you issue a `SELECT ... INTO OUTFILE` or `LOAD DATA INFILE` statement, your row in the `user` table does not have the `FILE` privilege enabled.
- If you change the grant tables directly (for example, by using `INSERT`, `UPDATE`, or `DELETE` statements) and your changes seem to be ignored, remember that you must execute a `FLUSH PRIVILEGES` statement or a `mysqladmin flush-privileges` command to cause the server to reload the privilege tables. Otherwise, your changes have no effect until the next time the server is restarted. Remember that after you change the `root` password with an `UPDATE` statement, you will not need to specify the new password until after you flush the privileges, because the server will not know you've changed the password yet!
- If your privileges seem to have changed in the middle of a session, it may be that a MySQL administrator has changed them. Reloading the grant tables affects new client connections, but it also affects existing connections as indicated in Section 6.2.8, “When Privilege Changes Take Effect”.
- If you have access problems with a Perl, PHP, Python, or ODBC program, try to connect to the server with `mysql -u user_name db_name` or `mysql -u user_name -pyour_pass db_name`. If you are able to connect using the `mysql` client, the problem lies with your program, not with the access privileges. (There is no space between `-p` and the password; you can also use the `--password=your_pass` syntax to specify the password. If you use the `-p` or `--password` option with no password value, MySQL prompts you for the password.)
- For testing purposes, start the `mysqld` server with the `--skip-grant-tables` option. Then you can change the MySQL grant tables and use the `SHOW GRANTS` statement to check whether your modifications have the desired effect. When you are satisfied with your changes, execute `mysqladmin flush-privileges` to tell the `mysqld` server to reload the privileges. This enables you to begin using the new grant table contents without stopping and restarting the server.
- If everything else fails, start the `mysqld` server with a debugging option (for example, `--debug=d,general,query`). This prints host and user information about attempted connections, as well as information about each command issued. See Section 28.5.3, “The DBUG Package”.
- If you have any other problems with the MySQL grant tables and feel you must post the problem to the mailing list, always provide a dump of the MySQL grant tables. You can dump the tables with the `mysqldump mysql` command. To file a bug report, see the instructions at Section 1.7, “How to Report Bugs or Problems”. In some cases, you may need to restart `mysqld` with `--skip-grant-tables` to run `mysqldump`.

6.3 MySQL User Account Management

This section describes how to set up accounts for clients of your MySQL server. It discusses the following topics:

- The meaning of account names and passwords as used in MySQL and how that compares to names and passwords used by your operating system
- How to set up new accounts and remove existing accounts
- How to use roles, which are named collections of privileges
- How to change passwords
- Guidelines for using passwords securely

See also Section 13.7.1, “Account Management Statements”, which describes the syntax and use for all user-management SQL statements.

6.3.1 User Names and Passwords

MySQL stores accounts in the `user` table of the `mysql` system database. An account is defined in terms of a user name and the client host or hosts from which the user can connect to the server. For information about account representation in the `user` table, see Section 6.2.3, “Grant Tables”.

The account may also have a password. MySQL supports authentication plugins, so it is possible that an account authenticates using some external authentication method. See Section 6.3.10, “Pluggable Authentication”.

There are several distinctions between the way user names and passwords are used by MySQL and your operating system:

- User names, as used by MySQL for authentication purposes, have nothing to do with user names (login names) as used by Windows or Unix. On Unix, most MySQL clients by default try to log in using the current Unix user name as the MySQL user name, but that is for convenience only. The default can be overridden easily, because client programs permit any user name to be specified with a `-u` or `--user` option. This means that anyone can attempt to connect to the server using any user name, so you cannot make a database secure in any way unless all MySQL accounts have passwords. Anyone who specifies a user name for an account that has no password is able to connect successfully to the server.

- MySQL user names can be up to 32 characters long. Operating system user names may be of a different maximum length. For example, Unix user names typically are limited to eight characters.

  Warning

  The limit on MySQL user name length is hardcoded in MySQL servers and clients, and trying to circumvent it by modifying the definitions of the tables in the `mysql` database *does not work*.

  You should never alter the structure of tables in the `mysql` database in any manner whatsoever except by means of the procedure that is described in Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”. Attempting to redefine MySQL's system tables in any other fashion results in undefined (and unsupported!) behavior. The server is free to ignore rows that become malformed as a result of such modifications.

- To authenticate client connections for accounts that use MySQL native authentication (implemented by the `mysql_native_password` authentication plugin), the server uses passwords stored in the `user` table. These passwords are distinct from passwords for logging in to your operating system. There is no necessary connection between the “external” password you use to log in to a Windows or Unix machine and the password you use to access the MySQL server on that machine.

  If the server authenticates a client using some other plugin, the authentication method that the plugin implements may or may not use a password stored in the `user` table. In this case, it is possible that an external password is also used to authenticate to the MySQL server.

- Passwords stored in the `user` table are encrypted using plugin-specific algorithms.

- If the user name and password contain only ASCII characters, it is possible to connect to the server regardless of character set settings. To connect when the user name or password contain non-ASCII characters, the client should call the `mysql_options()` C API function with the `MYSQL_SET_CHARSET_NAME` option and appropriate character set name as arguments. This causes authentication to take place using the specified character set. Otherwise, authentication will fail unless the server default character set is the same as the encoding in the authentication defaults.

  Standard MySQL client programs support a `--default-character-set` option that causes `mysql_options()` to be called as just described. In addition, character set autodetection is supported as described in Section 10.1.4, “Connection Character Sets and Collations”. For programs

that use a connector that is not based on the C API, the connector may provide an equivalent to `mysql_options()` that can be used instead. Check the connector documentation.

The preceding notes do not apply for `ucs2`, `utf16`, and `utf32`, which are not permitted as client character sets.

The MySQL installation process populates the grant tables with an initial `root` account, as described in Section 2.9.4, “Securing the Initial MySQL Account”, which also discusses how to assign passwords to it. Thereafter, you normally set up, modify, and remove MySQL accounts using statements such as `CREATE USER`, `DROP USER`, `GRANT`, and `REVOKE`. See Section 13.7.1, “Account Management Statements”.

To connect to a MySQL server with a command-line client, specify user name and password options as necessary for the account that you want to use:

```
shell> mysql --user=monty --password db_name
```

If you prefer short options, the command looks like this:

```
shell> mysql -u monty -p db_name
```

If you omit the password value following the `--password` or `-p` option on the command line (as just shown), the client prompts for one. Alternatively, the password can be specified on the command line:

```
shell> mysql --user=monty --password=password db_name
shell> mysql -u monty -ppassword db_name
```

If you use the `-p` option, there must be *no space* between `-p` and the following password value.

Specifying a password on the command line should be considered insecure. See Section 6.1.2.1, “End-User Guidelines for Password Security”. You can use an option file or a login path file to avoid giving the password on the command line. See Section 4.2.6, “Using Option Files”, and Section 4.6.7, “`mysql_config_editor` — MySQL Configuration Utility”.

For additional information about specifying user names, passwords, and other connection parameters, see Section 4.2.2, “Connecting to the MySQL Server”.

6.3.2 Adding User Accounts

You can create MySQL accounts two ways:

- By using account-management statements intended for creating accounts and establishing their privileges, such as `CREATE USER` and `GRANT`. These statements cause the server to make appropriate modifications to the underlying grant tables.
- By manipulating the MySQL grant tables directly with statements such as `INSERT`, `UPDATE`, or `DELETE`.

The preferred method is to use account-management statements because they are more concise and less error-prone than manipulating the grant tables directly. All such statements are described in Section 13.7.1, “Account Management Statements”. Direct grant table manipulation is discouraged, and is not described here. The server is free to ignore rows that become malformed as a result of such modifications.

Another option for creating accounts is to use the GUI tool MySQL Workbench. Also, several third-party programs offer capabilities for MySQL account administration. `phpMyAdmin` is one such program.

The following examples show how to use the `mysql` client program to set up new accounts. These examples assume that privileges have been set up according to the defaults described in Section 2.9.4,

"Securing the Initial MySQL Account". This means that to make changes, you must connect to the MySQL server as the MySQL `root` user, which has the `CREATE USER` privilege.

First, use the `mysql` program to connect to the server as the MySQL `root` user:

```
shell> mysql --user=root mysql
```

If you have assigned a password to the `root` account, you must also supply a `--password` or `-p` option.

After connecting to the server as `root`, you can add new accounts. The following example uses `CREATE USER` and `GRANT` statements to set up four accounts:

```
mysql> CREATE USER 'monty'@'localhost' IDENTIFIED BY 'some_pass';
mysql> GRANT ALL PRIVILEGES ON *.* TO 'monty'@'localhost'
    ->     WITH GRANT OPTION;
mysql> CREATE USER 'monty'@'%' IDENTIFIED BY 'some_pass';
mysql> GRANT ALL PRIVILEGES ON *.* TO 'monty'@'%'
    ->     WITH GRANT OPTION;
mysql> CREATE USER 'admin'@'localhost' IDENTIFIED BY 'admin_pass';
mysql> GRANT RELOAD,PROCESS ON *.* TO 'admin'@'localhost';
mysql> CREATE USER 'dummy'@'localhost';
```

The accounts created by those statements have the following properties:

- Two accounts have a user name of `monty` and a password of `some_pass`. Both are superuser accounts with full privileges to do anything. The `'monty'@'localhost'` account can be used only when connecting from the local host. The `'monty'@'%'` account uses the `'%'` wildcard for the host part, so it can be used to connect from any host.

  The `'monty'@'localhost'` account is necessary if there is an anonymous-user account for `localhost`. Without the `'monty'@'localhost'` account, that anonymous-user account takes precedence when `monty` connects from the local host and `monty` is treated as an anonymous user. The reason for this is that the anonymous-user account has a more specific `Host` column value than the `'monty'@'%'` account and thus comes earlier in the `user` table sort order. (`user` table sorting is discussed in Section 6.2.6, "Access Control, Stage 1: Connection Verification".)

- The `'admin'@'localhost'` account has a password of `admin_pass`. This account can be used only by `admin` to connect from the local host. It is granted the `RELOAD` and `PROCESS` administrative privileges. These privileges enable the `admin` user to execute the `mysqladmin reload`, `mysqladmin refresh`, and `mysqladmin flush-xxx` commands, as well as `mysqladmin processlist` . No privileges are granted for accessing any databases. You could add such privileges using `GRANT` statements.

- The `'dummy'@'localhost'` account has no password (which is insecure and not recommended). This account can be used only to connect from the local host. No privileges are granted. It is assumed that you will grant specific privileges to the account using `GRANT` statements.

To see the privileges for an account, use `SHOW GRANTS`:

```
mysql> SHOW GRANTS FOR 'admin'@'localhost';
+-----------------------------------------------------+
| Grants for admin@localhost                          |
+-----------------------------------------------------+
| GRANT RELOAD, PROCESS ON *.* TO 'admin'@'localhost' |
+-----------------------------------------------------+
```

To see nonprivilege properties for an account, use `SHOW CREATE USER`:

```
mysql> SHOW CREATE USER 'admin'@'localhost'\G
```

```
*************************** 1. row ***************************
CREATE USER for admin@localhost: CREATE USER 'admin'@'localhost'
IDENTIFIED WITH 'mysql_native_password'
AS '*67ACDEBDAB923990001F0FFB017EB8ED41861105'
REQUIRE NONE PASSWORD EXPIRE DEFAULT ACCOUNT UNLOCK
```

The next examples create three accounts and grant them access to specific databases. Each of them has a user name of `custom` and password of `obscure`:

```
mysql> CREATE USER 'custom'@'localhost' IDENTIFIED BY 'obscure';
mysql> GRANT SELECT,INSERT,UPDATE,DELETE,CREATE,DROP
    ->     ON bankaccount.*
    ->     TO 'custom'@'localhost';
mysql> CREATE USER 'custom'@'host47.example.com' IDENTIFIED BY 'obscure';
mysql> GRANT SELECT,INSERT,UPDATE,DELETE,CREATE,DROP
    ->     ON expenses.*
    ->     TO 'custom'@'host47.example.com';
mysql> CREATE USER 'custom'@'%.example.com' IDENTIFIED BY 'obscure';
mysql> GRANT SELECT,INSERT,UPDATE,DELETE,CREATE,DROP
    ->     ON customer.*
    ->     TO 'custom'@'%.example.com';
```

The three accounts can be used as follows:

- The first account can access the `bankaccount` database, but only from the local host.
- The second account can access the `expenses` database, but only from the host `host47.example.com`.
- The third account can access the `customer` database, from any host in the `example.com` domain. This account has access from all machines in the domain due to use of the `%` wildcard character in the host part of the account name.

6.3.3 Removing User Accounts

To remove an account, use the `DROP USER` statement, which is described in Section 13.7.1.5, “DROP USER Syntax”. For example:

```
mysql> DROP USER 'jeffrey'@'localhost';
```

6.3.4 Using Roles

A MySQL role is a named collection of privileges. Like user accounts, roles can have privileges granted to and revoked from them.

A user account can be granted roles, which grants to the account the role privileges. This enables assignment of sets of privileges to accounts and provides a convenient alternative to granting individual privileges, both for conceptualizing desired privilege assignments and implementing them.

The following list summarizes role management capabilities provided by MySQL:

- `CREATE ROLE` and `DROP ROLE` enable roles to be created and removed.
- `GRANT` and `REVOKE` enable privilege assignment and revocation for user accounts and roles.
- `SHOW GRANTS` displays privilege and role assignments for user accounts and roles.
- `SET DEFAULT ROLE` specifies which account roles are active by default.
- `SET ROLE` changes the active roles within the current session.
- The `CURRENT_ROLE()` function displays the active roles within the current session.

For descriptions of individual role-manipulation statements, see Section 13.7.1, “Account Management Statements”. The following discussion provides examples of role usage. Unless otherwise specified, SQL statements shown here should be executed using an account with administrative privileges, such as the `root` account.

Consider this scenario:

- An application uses a database named `app_db`.
- Associated with the application, there can be accounts for developers who create and maintain the application, and for users who interact with it.
- Developers need full access to the database. Some users need only read access, others need read/write access.

To avoid granting privileges individually to possibly many user accounts, create roles as names for the required privilege sets. This makes it easy to grant the required privileges to user accounts, by granting the appropriate roles.

To create the roles, use `CREATE ROLE`:

```
CREATE ROLE 'app_developer', 'app_read', 'app_write';
```

Role names are much like user account names and consist of a user part and host part in *`'user_name'@'host_name'`* format. The host part, if omitted, defaults to `'%'`. The user and host parts can be unquoted unless they contain special characters such as `-` or `%`. Unlike account names, the user part of role names cannot be blank. For additional information, see Section 6.2.5, “Specifying Role Names”.

To assign privileges to the roles, execute `GRANT` using the same syntax as for assigning privileges to user accounts:

```
GRANT ALL ON app_db.* TO 'app_developer';
GRANT SELECT ON app_db.* TO 'app_read';
GRANT INSERT, UPDATE, DELETE ON app_db.* TO 'app_write';
```

Now suppose that initially you require one developer account, two user accounts that need read-only access, and one user account that needs read/write access. Use `CREATE USER` to create the accounts:

```
CREATE USER 'dev1'@'localhost' IDENTIFIED BY 'dev1pass';
CREATE USER 'read_user1'@'localhost' IDENTIFIED BY 'read_user1pass';
CREATE USER 'read_user2'@'localhost' IDENTIFIED BY 'read_user2pass';
CREATE USER 'rw_user1'@'localhost' IDENTIFIED BY 'rw_user1pass';
```

To assign each user account its required privileges, you could use `GRANT` statements of the same form as just shown, but that requires enumerating individual privileges for each user. Instead, use an alternative `GRANT` syntax that permits granting roles rather than privileges:

```
GRANT 'app_developer' TO 'dev1'@'localhost';
GRANT 'app_read' TO 'read_user1'@'localhost', 'read_user2'@'localhost';
GRANT 'app_read', 'app_write' TO 'rw_user1'@'localhost';
```

The `GRANT` statement for the `rw_user1` account grants the read and write roles, which combine to provide the required read and write privileges.

The `GRANT` syntax for granting roles to an account differs from the syntax for granting privileges: There is an `ON` clause to assign privileges, whereas there is no `ON` clause to assign roles. Because

the syntaxes are distinct, you cannot mix assigning privileges and roles in the same statement. (It is permitted to assign both privileges and roles to an account, but you must use separate `GRANT` statements, each with syntax appropriate to what is to be granted.)

To verify the privilege assignments, use `SHOW GRANTS`. For example:

```
mysql> SHOW GRANTS FOR 'dev1'@'localhost';
+-----------------------------------------------+
| Grants for dev1@localhost                     |
+-----------------------------------------------+
| GRANT USAGE ON *.* TO `dev1`@`localhost`      |
| GRANT `app_developer`@`%` TO `dev1`@`localhost` |
+-----------------------------------------------+
```

However, that shows each granted role without “expanding” it to the privileges the role represents. To show role privileges as well, add a `USING` clause naming the granted roles for which to display privileges:

```
mysql> SHOW GRANTS FOR 'dev1'@'localhost' USING 'app_developer';
+-----------------------------------------------------------+
| Grants for dev1@localhost                                 |
+-----------------------------------------------------------+
| GRANT USAGE ON *.* TO `dev1`@`localhost`                  |
| GRANT ALL PRIVILEGES ON `app_db`.* TO `dev1`@`localhost` |
| GRANT `app_developer`@`%` TO `dev1`@`localhost`           |
+-----------------------------------------------------------+
```

Unless you are displaying information for your own account, the `USING` clause requires the `SUPER` privilege.

Verify each other type of user similarly:

```
mysql> SHOW GRANTS FOR 'read_user1'@'localhost' USING 'app_read';
+-----------------------------------------------------+
| Grants for read_user1@localhost                     |
+-----------------------------------------------------+
| GRANT USAGE ON *.* TO `read_user1`@`localhost`      |
| GRANT SELECT ON `app_db`.* TO `read_user1`@`localhost` |
| GRANT `app_read`@`%` TO `read_user1`@`localhost`    |
+-----------------------------------------------------+
mysql> SHOW GRANTS FOR 'rw_user1'@'localhost' USING 'app_read', 'app_write';
+-------------------------------------------------------------------------------+
| Grants for rw_user1@localhost                                                 |
+-------------------------------------------------------------------------------+
| GRANT USAGE ON *.* TO `rw_user1`@`localhost`                                  |
| GRANT SELECT, INSERT, UPDATE, DELETE ON `app_db`.* TO `rw_user1`@`localhost` |
| GRANT `app_read`@`%`,`app_write`@`%` TO `rw_user1`@`localhost`                |
+-------------------------------------------------------------------------------+
```

Granting roles to a user account does not automatically cause the roles to become active within sessions for the account. For example, if you connect to the server as `rw_user1` and invoke the `CURRENT_ROLE()` function, the result is `NONE` (no active roles):

```
mysql> SELECT CURRENT_ROLE();
+----------------+
| CURRENT_ROLE() |
+----------------+
| NONE           |
+----------------+
```

To specify which roles should become active each time a user connects to the server and authenticates, use `SET DEFAULT ROLE`. To set the default for each account to all its assigned roles, use this statement:

```
SET DEFAULT ROLE ALL TO
  'dev1'@'localhost',
  'read_user1'@'localhost',
  'read_user2'@'localhost',
  'rw_user1'@'localhost';
```

Now if you connect as `rw_user1`, the initial value of `CURRENT_ROLE()` reflects the new default role assignments:

```
mysql> SELECT CURRENT_ROLE();
+-------------------------------+
| CURRENT_ROLE()                |
+-------------------------------+
| `app_read`@`%`,`app_write`@`%` |
+-------------------------------+
```

Within a session, a user can execute `SET ROLE` to change the set of active roles. For example, for `rw_user1`:

```
mysql> SET ROLE NONE; SELECT CURRENT_ROLE();
+----------------+
| CURRENT_ROLE() |
+----------------+
| NONE           |
+----------------+
mysql> SET ROLE ALL EXCEPT 'app_write'; SELECT CURRENT_ROLE();
+----------------+
| CURRENT_ROLE() |
+----------------+
| `app_read`@`%` |
+----------------+
mysql> SET ROLE DEFAULT; SELECT CURRENT_ROLE();
+-------------------------------+
| CURRENT_ROLE()                |
+-------------------------------+
| `app_read`@`%`,`app_write`@`%` |
+-------------------------------+
```

The first `SET ROLE` statement deactivates all roles. The second makes `rw_user1` effectively read only. The third restores the default roles.

The privileges assigned to a role can be modified. This affects not only the role itself, but any account granted that role. Suppose that you want to temporarily make all application users read only. To do this, use `REVOKE` to revoke the modification privileges from the `app_write` role:

```
REVOKE INSERT, UPDATE, DELETE ON app_db.* FROM 'app_write';
```

As it happens, that leaves the role with no privileges at all, as can be seen using `SHOW GRANTS` (which demonstrates that this statement can be used with roles, not just users):

```
mysql> SHOW GRANTS FOR 'app_write';
+---------------------------------------+
| Grants for app_write@%                |
+---------------------------------------+
| GRANT USAGE ON *.* TO `app_write`@`%` |
+---------------------------------------+
```

Because revoking privileges from a role affects the privileges for any user assigned the modified role, `rw_user1` now has no modification privileges (`INSERT`, `UPDATE`, and `DELETE` are no longer present):

```
mysql> SHOW GRANTS FOR 'rw_user1'@'localhost'
       USING 'app_read', 'app_write';
```

```
+-------------------------------------------------------------+
| Grants for rw_user1@localhost                               |
+-------------------------------------------------------------+
| GRANT USAGE ON *.* TO `rw_user1`@`localhost`                 |
| GRANT SELECT ON `app_db`.* TO `rw_user1`@`localhost`         |
| GRANT `app_read`@`%`,`app_write`@`%` TO `rw_user1`@`localhost` |
+-------------------------------------------------------------+
```

In effect, the `rw_user1` read/write user has become a read-only user. This also occurs for any other accounts that are granted the `app_write` role. This illustrates how use of roles makes it unnecessary to modify privileges for individual accounts.

To restore modification privileges to the role, simply re-grant them:

```
GRANT INSERT, UPDATE, DELETE ON app_db.* TO 'app_write';
```

Now `rw_user1` again has modification privileges, as do any other accounts granted the `app_write` role.

To remove roles, use `DROP ROLE`:

```
DROP ROLE 'app_read', 'app_write';
```

Dropping a role revokes it from every account to which it was granted.

As has been hinted at earlier for `SHOW GRANTS`, which displays grants for user accounts or roles, accounts and roles can be used interchangeably. You can treat a user account like a role and grant that account to another user or a role. The effect is to grant the account's privileges and roles to the other user or role.

This set of statements demonstrates that you can grant a user to a user, a role to a user, a user to a role, or a role to a role:

```
CREATE USER 'u1';
CREATE ROLE 'r1';
GRANT SELECT ON db1.* TO 'u1';
GRANT SELECT ON db2.* TO 'r1';
CREATE USER 'u2';
CREATE ROLE 'r2';
GRANT 'u1', 'r1' TO 'u2';
GRANT 'u1', 'r1' TO 'r2';
```

The result in each case is to grant to the grantee object the privileges associated with the granted object. After executing those statements, each of `u2` and `r2` have been granted privileges from a user (`u1`) and a role (`r1`):

```
mysql> SHOW GRANTS FOR 'u2' USING 'u1', 'r1';
+-----------------------------------+
| Grants for u2@%                   |
+-----------------------------------+
| GRANT USAGE ON *.* TO `u2`@`%`    |
| GRANT SELECT ON `db1`.* TO `u2`@`%` |
| GRANT SELECT ON `db2`.* TO `u2`@`%` |
| GRANT `u1`@`%`,`r1`@`%` TO `u2`@`%` |
+-----------------------------------+
mysql> SHOW GRANTS FOR 'r2' USING 'u1', 'r1';
+-----------------------------------+
| Grants for r2@%                   |
+-----------------------------------+
| GRANT USAGE ON *.* TO `r2`@`%`    |
| GRANT SELECT ON `db1`.* TO `r2`@`%` |
| GRANT SELECT ON `db2`.* TO `r2`@`%` |
```

```
| GRANT `u1`@`%`,`r1`@`%` TO `r2`@`%` |
+-------------------------------------+
```

The preceding example is illustrative only, but interchangeability of user accounts and roles has practical application, such as in the following situation: Suppose that a legacy application development project began before the advent of roles in MySQL, so all user accounts associated with the project are granted privileges directly (rather than granted privileges by virtue of being granted roles). One of these accounts is a developer account that was originally granted privileges as follows:

```
CREATE USER 'old_app_dev'@'localhost' IDENTIFIED BY 'old_app_devpass';
GRANT ALL ON old_app.* TO 'old_app_dev'@'localhost';
```

If this developer leaves the project, it becomes necessary to assign the privileges to another user, or perhaps multiple users if development activies have expanded. Here are some ways to deal with the issue:

- Without using roles: Change the account password so the original developer cannot use it, and have a new developer use the account instead:

  ```
  ALTER USER 'old_app_dev'@'localhost' IDENTIFIED BY 'new_password';
  ```

- Using roles: Lock the account to prevent anyone from using it to connect to the server:

  ```
  ALTER USER 'old_app_dev'@'localhost' ACCOUNT LOCK;
  ```

  Then treat the account as a role. For each developer new to the project, create a new account and grant to it the original developer account:

  ```
  CREATE USER 'new_app_dev1'@'localhost' IDENTIFIED BY 'new_password';
  GRANT 'old_app_dev'@'localhost' TO 'new_app_dev1'@'localhost';
  ```

  The effect is to assign the original developer account privileges to the new account.

6.3.5 Reserved User Accounts

One part of the MySQL installation process is data directory initialization (see Section 2.9.1.1, “Initializing the Data Directory Manually Using mysqld”). During data directory initialization, MySQL creates user accounts that should be considered reserved:

- `'root'@'localhost'`: Used for administrative purposes. This account has all privileges and can perform any operation.

  Strictly speaking, this account name is not reserved, in the sense that some installations rename the `root` account to something else to avoid exposing a highly privileged account with a well-known name.

- `'mysql.sys'@'localhost'`: Used as the `DEFINER` for `sys` schema objects. Use of the `mysql.sys` account avoids problems that occur if a DBA renames or removes the `root` account. This account is locked so that it cannot be used for client conections.

- `'mysql.session'@'localhost'`: Used internally by plugins to access the server. This account is locked so that it cannot be used for client conections.

6.3.6 Setting Account Resource Limits

One means of restricting client use of MySQL server resources is to set the global `max_user_connections` system variable to a nonzero value. This limits the number of simultaneous connections that can be made by any given account, but places no limits on what a client can do once

connected. In addition, setting `max_user_connections` does not enable management of individual accounts. Both types of control are of interest to MySQL administrators.

To address such concerns, MySQL permits limits for individual accounts on use of these server resources:

- The number of queries an account can issue per hour
- The number of updates an account can issue per hour
- The number of times an account can connect to the server per hour
- The number of simultaneous connections to the server by an account

Any statement that a client can issue counts against the query limit, unless its results are served from the query cache. Only statements that modify databases or tables count against the update limit.

An “account” in this context corresponds to a row in the `mysql.user` table. That is, a connection is assessed against the `User` and `Host` values in the `user` table row that applies to the connection. For example, an account `'usera'@'%.example.com'` corresponds to a row in the `user` table that has `User` and `Host` values of `usera` and `%.example.com`, to permit `usera` to connect from any host in the `example.com` domain. In this case, the server applies resource limits in this row collectively to all connections by `usera` from any host in the `example.com` domain because all such connections use the same account.

Before MySQL 5.0.3, an “account” was assessed against the actual host from which a user connects. This older method of accounting may be selected by starting the server with the `--old-style-user-limits` option. In this case, if `usera` connects simultaneously from `host1.example.com` and `host2.example.com`, the server applies the account resource limits separately to each connection. If `usera` connects again from `host1.example.com`, the server applies the limits for that connection together with the existing connection from that host.

To establish resource limits for an account at account-creation time, use the `CREATE USER` statement. To modify the limits for an existing account, use `ALTER USER`. Provide a `WITH` clause that names each resource to be limited. The default value for each limit is zero (no limit). For example, to create a new account that can access the `customer` database, but only in a limited fashion, issue these statements:

```
mysql> CREATE USER 'francis'@'localhost' IDENTIFIED BY 'frank'
    ->     WITH MAX_QUERIES_PER_HOUR 20
    ->          MAX_UPDATES_PER_HOUR 10
    ->          MAX_CONNECTIONS_PER_HOUR 5
    ->          MAX_USER_CONNECTIONS 2;
```

The limit types need not all be named in the `WITH` clause, but those named can be present in any order. The value for each per-hour limit should be an integer representing a count per hour. For `MAX_USER_CONNECTIONS`, the limit is an integer representing the maximum number of simultaneous connections by the account. If this limit is set to zero, the global `max_user_connections` system variable value determines the number of simultaneous connections. If `max_user_connections` is also zero, there is no limit for the account.

To modify limits for an existing account, use an `ALTER USER` statement. The following statement changes the query limit for `francis` to 100:

```
mysql> ALTER USER 'francis'@'localhost' WITH MAX_QUERIES_PER_HOUR 100;
```

The statement modifies only the limit value specified and leaves the account otherwise unchanged.

To remove a limit, set its value to zero. For example, to remove the limit on how many times per hour `francis` can connect, use this statement:

```
mysql> ALTER USER 'francis'@'localhost' WITH MAX_CONNECTIONS_PER_HOUR 0;
```

As mentioned previously, the simultaneous-connection limit for an account is determined from the MAX_USER_CONNECTIONS limit and the max_user_connections system variable. Suppose that the global max_user_connections value is 10 and three accounts have individual resource limits specified as follows:

```
ALTER USER 'user1'@'localhost' WITH MAX_USER_CONNECTIONS 0;
ALTER USER 'user2'@'localhost' WITH MAX_USER_CONNECTIONS 5;
ALTER USER 'user3'@'localhost' WITH MAX_USER_CONNECTIONS 20;
```

user1 has a connection limit of 10 (the global max_user_connections value) because it has a MAX_USER_CONNECTIONS limit of zero. user2 and user3 have connection limits of 5 and 20, respectively, because they have nonzero MAX_USER_CONNECTIONS limits.

The server stores resource limits for an account in the user table row corresponding to the account. The max_questions, max_updates, and max_connections columns store the per-hour limits, and the max_user_connections column stores the MAX_USER_CONNECTIONS limit. (See Section 6.2.3, “Grant Tables”.)

Resource-use counting takes place when any account has a nonzero limit placed on its use of any of the resources.

As the server runs, it counts the number of times each account uses resources. If an account reaches its limit on number of connections within the last hour, the server rejects further connections for the account until that hour is up. Similarly, if the account reaches its limit on the number of queries or updates, the server rejects further queries or updates until the hour is up. In all such cases, the server issues appropriate error messages.

Resource counting occurs per account, not per client. For example, if your account has a query limit of 50, you cannot increase your limit to 100 by making two simultaneous client connections to the server. Queries issued on both connections are counted together.

The current per-hour resource-use counts can be reset globally for all accounts, or individually for a given account:

- To reset the current counts to zero for all accounts, issue a FLUSH USER_RESOURCES statement. The counts also can be reset by reloading the grant tables (for example, with a FLUSH PRIVILEGES statement or a mysqladmin reload command).
- The counts for an individual account can be reset to zero by setting any of its limits again. Specify a limit value equal to the value currently assigned to the account.

Per-hour counter resets do not affect the MAX_USER_CONNECTIONS limit.

All counts begin at zero when the server starts. Counts do not carry over through server restarts.

For the MAX_USER_CONNECTIONS limit, an edge case can occur if the account currently has open the maximum number of connections permitted to it: A disconnect followed quickly by a connect can result in an error (ER_TOO_MANY_USER_CONNECTIONS or ER_USER_LIMIT_REACHED) if the server has not fully processed the disconnect by the time the connect occurs. When the server finishes disconnect processing, another connection will once more be permitted.

6.3.7 Assigning Account Passwords

Required credentials for clients that connect to the MySQL server can include a password. This section describes how to assign passwords for MySQL accounts.

MySQL stores passwords in the user table in the mysql system database. Operations that assign or modify passwords are permitted only to users with the CREATE USER privilege, or, alternatively, privileges for the mysql database (INSERT privilege to create new accounts, UPDATE privilege to modify existing accounts). If the read_only system variable is enabled, use of account-modification

statements such as `CREATE USER` or `SET PASSWORD` additionally requires `CONNECTION_ADMIN` or the `SUPER` privilege.

The discussion here summarizes syntax only for the most common password-assignment statements. For complete details on other possibilities, see Section 13.7.1.3, “CREATE USER Syntax”, Section 13.7.1.1, “ALTER USER Syntax”, Section 13.7.1.6, “GRANT Syntax”, and Section 13.7.1.10, “SET PASSWORD Syntax”.

MySQL hashes passwords stored in the `mysql.user` table to obfuscate them. For most statements described here, MySQL automatically hashes the password specified. An exception is `SET PASSWORD ... = PASSWORD('auth_string')`, for which you use the `PASSWORD()` function explicitly to hash the password. There are also syntaxes for `CREATE USER`, `ALTER USER`, `GRANT`, and `SET PASSWORD` that permit hashed values to be specified literally; for details, see the descriptions of those statements.

MySQL uses plugins to perform client authentication; see Section 6.3.10, “Pluggable Authentication”. The authentication plugin associated with an account determines the algorithm used to hash passwords for that account.

To assign a password when you create a new account, use `CREATE USER` and include an `IDENTIFIED BY` clause:

```
mysql> CREATE USER 'jeffrey'@'localhost'
    -> IDENTIFIED BY 'mypass';
```

For this `CREATE USER` syntax, MySQL automatically hashes the password before storing it in the `mysql.user` table.

`CREATE USER` also supports syntax for specifying the account authentication plugin. See Section 13.7.1.3, “CREATE USER Syntax”.

To assign or change a password for an existing account, use one of the following methods:

- Use the `ALTER USER` statement with an `IDENTIFIED BY` clause:

  ```
  mysql> ALTER USER 'jeffrey'@'localhost'
      -> IDENTIFIED BY 'mypass';
  ```

  If you are not connected as an anonymous user, you can change your own password without naming your own account literally:

  ```
  mysql> ALTER USER USER()
      -> IDENTIFIED BY 'mypass';
  ```

  For these `ALTER USER` syntaxes, MySQL automatically hashes the password before storing it in the `mysql.user` table.

- Use `SET PASSWORD` with the `PASSWORD()` function:

  ```
  mysql> SET PASSWORD FOR
      -> 'jeffrey'@'localhost' = PASSWORD('mypass');
  ```

  If you are not connected as an anonymous user, you can change your own password by omitting the `FOR` clause:

  ```
  mysql> SET PASSWORD = PASSWORD('mypass');
  ```

  The `PASSWORD()` function hashes the password using the hashing method determined by the value of the `old_passwords` system variable value. If `SET PASSWORD` rejects the hashed password

value returned by `PASSWORD()` as not being in the correct format, it may be necessary to change `old_passwords` to change the hashing method. See Section 13.7.1.10, "SET PASSWORD Syntax".

Note

Use of `SET PASSWORD ... = PASSWORD('auth_string')` for password modification is deprecated as of MySQL 5.7.6. Use `ALTER USER` instead.

- Use `SET PASSWORD` without the `PASSWORD()` function:

  `SET PASSWORD` interprets the string as a cleartext string and hashes it appropriately for the account authentication plugin before storing it in the `mysql.user` account row.

```
mysql> SET PASSWORD FOR
    -> 'jeffrey'@'localhost' = 'mypass';
```

- Use a `GRANT USAGE` statement at the global level (`ON *.*`) to change an account password without affecting the account's current privileges:

```
mysql> GRANT USAGE ON *.* TO 'jeffrey'@'localhost'
    -> IDENTIFIED BY 'mypass';
```

  For this `GRANT` syntax, MySQL automatically hashes the password before storing it in the `mysql.user` table.

Note

Use of `GRANT` for password modification is deprecated as of MySQL 5.7.6. Use `ALTER USER` instead.

- To change an account password from the command line, use the `mysqladmin` command:

```
shell> mysqladmin -u user_name -h host_name password "new_password"
```

  The account for which this command sets the password is the one with a `mysql.user` table row that matches *`user_name`* in the `User` column and the client host *from which you connect* in the `Host` column.

  For password changes made using `mysqladmin`, MySQL automatically hashes the password before storing it in the `mysql.user` table.

6.3.8 Password Expiration Policy

MySQL enables database administrators to expire account passwords manually, and to establish a policy for automatic password expiration.

To expire a password manually, the database administrator uses the `ALTER USER` statement:

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE;
```

This operation marks the password expired in the corresponding `mysql.user` table row. The `mysql.user` table indicates for each account when its password was last changed, and the server automatically treats the password as expired at client connection time if it is past its permitted lifetime. This works with no explicit manual password expiration.

The `default_password_lifetime` system variable defines the global automatic password expiration policy. It applies to accounts that use MySQL built-in authentication methods (accounts that use an authentication plugin of `mysql_native_password` or `sha256_password`).

The default `default_password_lifetime` value is 0, which disables automatic password expiration. If the value of `default_password_lifetime` is a positive integer *`N`*, it indicates the permitted password lifetime; passwords must be changed every *`N`* days.

Examples:

- To establish a global policy that passwords have a lifetime of approximately six months, start the server with these lines in an option file:

```
[mysqld]
default_password_lifetime=180
```

- To establish a global policy such that passwords never expire, set `default_password_lifetime` to 0:

```
[mysqld]
default_password_lifetime=0
```

- `default_password_lifetime` can also be changed at runtime (this requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege):

```
SET GLOBAL default_password_lifetime = 180;
SET GLOBAL default_password_lifetime = 0;
```

No matter the global policy, it can be overridden for individual accounts with `ALTER USER`:

- Require the password to be changed every 90 days:

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE INTERVAL 90 DAY;
```

- Disable password expiration:

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE NEVER;
```

- Defer to the global expiration policy:

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE DEFAULT;
```

  `PASSWORD EXPIRE DEFAULT` defers to the global exipration policy and in the `mysql.user` table sets the `password_lifetime` field to `NULL` for the named account.

These `ALTER USER` statements update the corresponding `mysql.user` table row.

When a client successfully connects, the server determines whether the account password is expired:

- The server checks whether the password has been manually expired and, if so, restricts the session.
- Otherwise, the server checks whether the password is past its lifetime according to the automatic password expiration policy. If so, the server considers the password expired and restricts the session.

A restricted client operates in “sandbox mode,”, which limits the operations permitted to the client (see Section 6.3.9, “Password Expiration and Sandbox Mode”). Operations performed by a restricted client result in an error until the user establishes a new account password:

```
mysql> SELECT 1;
ERROR 1820 (HY000): You must SET PASSWORD before executing this statement

mysql> ALTER USER USER() IDENTIFIED BY 'new_password';
```

```
Query OK, 0 rows affected (0.01 sec)

mysql> SELECT 1;
+---+
| 1 |
+---+
| 1 |
+---+
1 row in set (0.00 sec)
```

This restricted mode of operation permits `SET` statements, which is useful if the deprecated `SET PASSWORD` is used instead of `ALTER USER` and the account password has a hashing format that requires `old_passwords` to be set to a value different from its default.

It is possible for an administrative user to reset the account password, but any existing sessions for that account remain restricted. A client using the account must disconnect and reconnect before statements can be executed successfully.

Note

It is possible to “reset” a password by setting it to its current value. As a matter of good policy, it is preferable to choose a different password.

6.3.9 Password Expiration and Sandbox Mode

MySQL provides password-expiration capability to enable database administrators to expire account passwords and require users to reset their password. This section describes how password expiration works.

To expire an account password, use the `ALTER USER` statement. For example:

```
ALTER USER 'myuser'@'localhost' PASSWORD EXPIRE;
```

This statement modifies the row of the `mysql.user` table associated with the named account, setting the `password_expired` column to `'Y'`. This does not affect any current connections the account has open. For each subsequent connection that uses the account, the server either disconnects the client or handles the client in “sandbox mode,” in which the server permits to the client only those operations necessary to reset the expired password. The action taken by the server depends on both client and server settings.

If the server disconnects the client, it returns an `ER_MUST_CHANGE_PASSWORD_LOGIN` error:

```
shell> mysql -u myuser -p
Password: ******
ERROR 1862 (HY000): Your password has expired. To log in you must
change it using a client that supports expired passwords.
```

If the server puts the client in sandbox mode, these operations are permitted within the client session:

- The client can reset the account password with `ALTER USER` or `SET PASSWORD`. This modifies the row of the `mysql.user` table associated with the current account, setting the `password_expired` column to `'N'`. After the password has been reset, the server restores normal access for the session, as well as for subsequent connections that use the account.

  It is possible to “reset” a password by setting it to its current value. As a matter of good policy, it is preferable to choose a different password.

- The client can use `SET` statements, which is useful if the deprecated `SET PASSWORD` is used instead of `ALTER USER` and the account password has a hashing format that requires `old_passwords` to be set to a value different from its default.

For any operation not permitted within the session, the server returns an `ER_MUST_CHANGE_PASSWORD` error:

```
mysql> USE performance_schema;
ERROR 1820 (HY000): You must SET PASSWORD before executing this statement

mysql> SELECT 1;
ERROR 1820 (HY000): You must SET PASSWORD before executing this statement
```

For noninteractive invocations of the `mysql` client (for example, in batch mode), the server normally disconnects the client if the password is expired. To permit `mysql` to stay connected so that the password can be changed (using the statements just described), add the `--connect-expired-password` option to the `mysql` command.

As mentioned previously, whether the server disconnects an expired-password client or puts it in sandbox mode depends on a combination of client and server settings. The following discussion describes the relevant settings and how they interact.

On the client side, a given client indicates whether it can handle sandbox mode for expired passwords. For clients that use the C client library, there are two ways to do this:

- Pass the `MYSQL_OPT_CAN_HANDLE_EXPIRED_PASSWORDS` flag to `mysql_options()` prior to connecting:

  ```
  arg = 1;
  result = mysql_options(mysql,
                         MYSQL_OPT_CAN_HANDLE_EXPIRED_PASSWORDS, &arg);
  ```

  The `mysql` client enables `MYSQL_OPT_CAN_HANDLE_EXPIRED_PASSWORDS` if invoked interactively or the `--connect-expired-password` option is given.

- Pass the `CLIENT_CAN_HANDLE_EXPIRED_PASSWORDS` flag to `mysql_real_connect()` at connection time:

  ```
  mysql = mysql_real_connect(mysql,
                             host, user, password, "test",
                             port, unix_socket,
                             CLIENT_CAN_HANDLE_EXPIRED_PASSWORDS);
  ```

Other MySQL Connectors have their own conventions for indicating readiness to handle sandbox mode. See the relevant Connector documentation.

On the server side, if a client indicates that it can handle expired passwords, the server puts it in sandbox mode.

If a client does not indicate that it can handle expired passwords (or uses an older version of the client library that cannot so indicate), the server action depends on the value of the `disconnect_on_expired_password` system variable:

- If `disconnect_on_expired_password` is enabled (the default), the server disconnects the client with an `ER_MUST_CHANGE_PASSWORD_LOGIN` error.
- If `disconnect_on_expired_password` is disabled, the server puts the client in sandbox mode.

The preceding client and server settings apply only for accounts with expired passwords. If a client connects using a nonexpired password, the server handles the client normally.

6.3.10 Pluggable Authentication

When a client connects to the MySQL server, the server uses the user name provided by the client and the client host to select the appropriate account row from the `mysql.user` table. The server then

authenticates the client, determining from the account row which authentication plugin applies to the client. The server invokes that plugin to authenticate the user, and the plugin returns a status to the server indicating whether the user is permitted to connect. If the server cannot find the plugin, an error occurs and the connection attempt is rejected.

Pluggable authentication enables two important capabilities:

- **External authentication:** Pluggable authentication makes it possible for clients to connect to the MySQL server with credentials that are appropriate for authentication methods that store credentials elsewhere than in the `mysql.user` table. For example, plugins can be created to use external authentication methods such as PAM, Windows login IDs, LDAP, or Kerberos.
- **Proxy users:** If a user is permitted to connect, an authentication plugin can return to the server a user name different from the name of the connecting user, to indicate that the connecting user is a proxy for another user (the proxied user). While the connection lasts, the proxy user is treated, for purposes of access control, as having the privileges of the proxied user. In effect, one user impersonates another. For more information, see Section 6.3.11, “Proxy Users”.

Several authentication plugins are available in MySQL:

- A plugin that performs native authentication; that is, authentication based on the password hashing method in use from before the introduction of pluggable authentication in MySQL. The `mysql_native_password` plugin implements authentication based on this native password hashing method. See Section 6.5.1.1, “Native Pluggable Authentication”. Native authentication using `mysql_native_password` is the default for new accounts, unless the `default_authentication_plugin` system variable is set otherwise.
- A plugin that performs authentication using SHA-256 password hashing. This encryption is stronger than that available with native authentication. See Section 6.5.1.2, “SHA-256 Pluggable Authentication”.
- A client-side plugin that sends the password to the server without hashing or encryption. This plugin is used in conjunction with server-side plugins that require access to the password exactly as provided by the client user. See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.
- A plugin prevents all client connections to any account that uses it. Use cases for this plugin include accounts that must be able to execute stored programs and views with elevated privileges without exposing those privileges to ordinary users, and proxied accounts that should never permit direct login but are accessed only through proxy accounts. See Section 6.5.1.4, “No-Login Pluggable Authentication”.
- A plugin that authenticates clients that connect from the local host through the Unix socket file. See Section 6.5.1.5, “Socket Peer-Credential Pluggable Authentication”.
- A test plugin that checks account credentials and logs success or failure to the server error log. This plugin is intended for testing and development purposes, and as an example of how to write an authentication plugin. See Section 6.5.1.6, “Test Pluggable Authentication”.

Note

For information about current restrictions on the use of pluggable authentication, including which connectors support which plugins, see Section C.9, “Restrictions on Pluggable Authentication”.

Third-party connector developers should read that section to determine the extent to which a connector can take advantage of pluggable authentication capabilities and what steps to take to become more compliant.

If you are interested in writing your own authentication plugins, see Section 28.2.4.9, “Writing Authentication Plugins”.

Authentication Plugin Usage Instructions

This section provides general instructions for installing and using authentication plugins. For instructions specific to a given plugin, see the section that describes that plugin.

In general, pluggable authentication uses corresponding plugins on the server and client sides, so you use a given authentication method like this:

- If necessary, install the plugin library or libraries containing the appropriate plugins. On the server host, install the library containing the server-side plugin, so that the server can use it to authenticate client connections. Similarly, on each client host, install the library containing the client-side plugin for use by client programs. Authentication plugins that are built in need not be installed.
- For each MySQL account that you create, specify the appropriate server-side plugin to use for authentication.
- When a client connects, the server-side plugin tells the client program which client-side plugin to use for authentication.

In the case that an account uses an authentication method that is the default for both the server and the client program, the server need not communicate to the client which client-side plugin to use, and a round trip in client/server negotiation can be avoided. This is true for accounts that use native MySQL authentication (`mysql_native_password`).

The `--default-auth=`*`plugin_name`* option can be specified on the `mysql` command line as a hint about which client-side plugin the program can expect to use, although the server will override this if the server-side plugin associated with the user account requires a different client-side plugin.

If the client program does not find the client-side plugin, specify a `--plugin-dir=`*`dir_name`* option to indicate where the plugin is located.

Note

If you start the server with the `--skip-grant-tables` option, authentication plugins are not used even if loaded because the server performs no client authentication and permits any client to connect. Because this is insecure, you might want to use `--skip-grant-tables` in conjunction with `--skip-networking` to prevent remote clients from connecting.

6.3.11 Proxy Users

The MySQL server authenticates client connections by means of authentication plugins. The plugin that authenticates a given connection may request that the connecting (external) user be treated as a different user for privilege-checking purposes. This enables the external user to be a proxy for the second user; that is, to have the privileges of the second user:

- The external user is a “proxy user” (a user who can impersonate or become known as another user).
- The second user is a “proxied user” (a user whose identity and privileges can be taken on by a proxy user).

This section describes how the proxy user capability works. For general information about authentication plugins, see Section 6.3.10, “Pluggable Authentication”. For information about specific plugins, see Section 6.5.1, “Authentication Plugins”. For information about writing authentication plugins that support proxy users, see Implementing Proxy User Support in Authentication Plugins.

- Requirements for Proxy User Support
- Granting the Proxy Privilege

- Default Proxy Users
- Default Proxy User and Anonymous User Conflicts
- Server Support for Proxy User Mapping
- Proxy User System Variables

Requirements for Proxy User Support

For proxying to occur for a given authentication plugin, these conditions must be satisfied:

- Proxying must be supported, either by the plugin itself, or by the MySQL server on behalf of the plugin. In the latter case, server support may need to be enabled explicitly; see Server Support for Proxy User Mapping.
- The proxy user account must be set up to be authenticated by the plugin. Use the `CREATE USER` statement to associate an account with an authentication plugin, or `ALTER USER` to change its plugin.
- The proxied user account must be created and granted the privileges to be permitted for the proxy user. Use the `CREATE USER` and `GRANT` statements for this.
- The proxy user account must have the `PROXY` privilege for the proxied account. Use the `GRANT` statement for this.
- For a client connecting to the proxy account to be treated as a proxy user, an authentication plugin that supports proxying must return a user name different from the client user name, to indicate the user name for the proxied account that defines the privileges to be used for the proxy user.

  Alternatively, for plugins that are provided proxy mapping by the server, the proxied user is determined from the `PROXY` privilege held by the proxy user.

The proxy mechanism permits mapping only the client user name to the proxied user name. There is no provision for mapping host names. When a connecting client matches a proxy account, the server attempts to find a match for a proxied account using the user name returned by the authentication plugin and the host name of the proxy account.

Consider the following definitions:

```
-- create proxy user
CREATE USER 'employee_ext'@'localhost'
  IDENTIFIED WITH my_auth_plugin AS 'my_auth_string';

-- create proxied user and grant its privileges
CREATE USER 'employee'@'localhost'
  IDENTIFIED BY 'employee_pass';
GRANT ALL ON employees.*
  TO 'employee'@'localhost';

-- grant PROXY privilege for proxy user to proxied user
GRANT PROXY
  ON 'employee'@'localhost'
  TO 'employee_ext'@'localhost';
```

When a client connects as `employee_ext` from the local host, MySQL uses the plugin named `my_auth_plugin` to perform authentication. Suppose that `my_auth_plugin` returns a user name of `employee` to the server, based on the content of `'my_auth_string'` and perhaps by consulting some external authentication system. The name `employee` differs from `employee_ext`, so returning `employee` serves as a request to the server to treat the `employee_ext` client, for purposes of privilege checking, as the `employee` local user.

In this case, `employee_ext` is the proxy user and `employee` is the proxied user.

The server verifies that proxy authentication for `employee` is possible for the `employee_ext` user by checking whether `employee_ext` (the proxy user) has the `PROXY` privilege for `employee` (the proxied user). If this privilege has not been granted, an error occurs. Otherwise, the server checks statements executed during the client session by `employee_ext` against the privileges granted to `employee`. In this case, `employee_ext` can access tables in the `employees` database.

When proxying occurs, the `USER()` and `CURRENT_USER()` functions can be used to see the difference between the connecting user (the proxy user) and the account whose privileges apply during the current session (the proxied user). For the example just described, those functions return these values:

```
mysql> SELECT USER(), CURRENT_USER();
+------------------------+--------------------+
| USER()                 | CURRENT_USER()     |
+------------------------+--------------------+
| employee_ext@localhost | employee@localhost |
+------------------------+--------------------+
```

In the `CREATE USER` statement that creates the proxy user account, the `IDENTIFIED WITH` clause that names the authentication plugin is optionally followed by an `AS 'auth_string'` clause specifying a string that the server passes to the plugin when the user connects. If present, the string provides information that helps the plugin determine how to map the external client user name to a proxied user name. It is up to each plugin whether it requires the `AS` clause. If so, the format of the authentication string depends on how the plugin intends to use it. Consult the documentation for a given plugin for information about the authentication string values it accepts.

Granting the Proxy Privilege

The `PROXY` privilege is needed to enable an external user to connect as and have the privileges of another user. To grant this privilege, use the `GRANT` statement. For example:

```
GRANT PROXY ON 'proxied_user' TO 'proxy_user';
```

The statement creates a row in the `mysql.proxies_priv` grant table.

At connection time, *`proxy_user`* must represent a valid externally authenticated MySQL user, and *`proxied_user`* must represent a valid locally authenticated user. Otherwise, the connection attempt fails.

The corresponding `REVOKE` syntax is:

```
REVOKE PROXY ON 'proxied_user' FROM 'proxy_user';
```

MySQL `GRANT` and `REVOKE` syntax extensions work as usual. For example:

```
GRANT PROXY ON 'a' TO 'b', 'c', 'd';
GRANT PROXY ON 'a' TO 'd' WITH GRANT OPTION;
GRANT PROXY ON 'a' TO ''@'';
REVOKE PROXY ON 'a' FROM 'b', 'c', 'd';
```

In the preceding example, `''@''` is the default proxy user and means “any user.” Default proxy users are discussed in Default Proxy Users.

The `PROXY` privilege can be granted in these cases:

- By a user that has `GRANT PROXY ... WITH GRANT OPTION` for *`proxied_user`*.
- By *`proxied_user`* for itself: The value of `USER()` must exactly match `CURRENT_USER()` and *`proxied_user`*, for both the user name and host name parts of the account name.

The initial `root` account created during MySQL installation has the `PROXY ... WITH GRANT OPTION` privilege for `''@''`, that is, for all users and all hosts. This enables `root` to set up proxy users, as well as to delegate to other accounts the authority to set up proxy users. For example, `root` can do this:

```
CREATE USER 'admin'@'localhost' IDENTIFIED BY 'test';
GRANT PROXY ON ''@'' TO 'admin'@'localhost' WITH GRANT OPTION;
```

Those statements create an `admin` user that can manage all `GRANT PROXY` mappings. For example, `admin` can do this:

```
GRANT PROXY ON sally TO joe;
```

Default Proxy Users

To specify that some or all users should connect using a given authentication plugin, create a “blank” MySQL user, associate it with that plugin, and let the plugin return the real authenticated user name (if different from the blank user). For example, suppose that there exists a plugin named `ldap_auth` that implements LDAP authentication and maps connecting users onto either a developer or manager account. To set up proxying of users onto these accounts, use the following statements:

```
-- create default proxy user
CREATE USER ''@'' IDENTIFIED WITH ldap_auth AS 'O=Oracle, OU=MySQL';

-- create proxied users
CREATE USER 'developer'@'localhost' IDENTIFIED BY 'developer_pass';
CREATE USER 'manager'@'localhost' IDENTIFIED BY 'manager_pass';

-- grant PROXY privilege for default proxy user to proxied users
GRANT PROXY ON 'manager'@'localhost' TO ''@'';
GRANT PROXY ON 'developer'@'localhost' TO ''@'';
```

Now assume that a client connects as follows:

```
shell> mysql --user=myuser --password ...
Enter password: myuser_pass
```

The server will not find `myuser` defined as a MySQL user. But because there is a blank user account (`''@''`) that matches the client user name and host name, the server authenticates the client against that account: The server invokes the `ldap_auth` authentication plugin and passes `myuser` and `myuser_pass` to it as the user name and password.

If the `ldap_auth` plugin finds in the LDAP directory that `myuser_pass` is not the correct password for `myuser`, authentication fails and the server rejects the connection.

If the password is correct and `ldap_auth` finds that `myuser` is a developer, it returns the user name `developer` to the MySQL server, rather than `myuser`. Returning a user name different from the client user name of `myuser` signals to the server that it should treat `myuser` as a proxy. The server verifies that `''@''` can authenticate as `developer` (because it has the `PROXY` privilege to do so) and accepts the connection. The session proceeds with `myuser` having the privileges of `developer`, the proxied user. (These privileges should be set up by the DBA using `GRANT` statements, not shown.) The `USER()` and `CURRENT_USER()` functions return these values:

```
mysql> SELECT USER(), CURRENT_USER();
+------------------+--------------------+
| USER()           | CURRENT_USER()     |
+------------------+--------------------+
| myuser@localhost | developer@localhost |
+------------------+--------------------+
```

If the plugin instead finds in the LDAP directory that `myuser` is a manager, it returns `manager` as the user name and the session proceeds with `myuser` having the privileges of `manager`.

```
mysql> SELECT USER(), CURRENT_USER();
+------------------+------------------+
| USER()           | CURRENT_USER()   |
+------------------+------------------+
| myuser@localhost | manager@localhost |
+------------------+------------------+
```

For simplicity, external authentication cannot be multilevel: Neither the credentials for `developer` nor those for `manager` are taken into account in the preceding example. However, they are still used if a client tries to connect and authenticate directly as the `developer` or `manager` account, which is why those accounts should be assigned passwords.

Default Proxy User and Anonymous User Conflicts

If you intend to create a default proxy user, check for other existing “match any user” accounts that take precedence over the default proxy user because they can prevent that user from working as intended.

In the preceding discussion, the default proxy user account has `''` in the host part, which matches any host. If you set up a default proxy user, take care to also check whether nonproxy accounts exist with the same user part and `'%'` in the host part, because `'%'` also matches any host, but has precedence over `''` by the rules that the server uses to sort account rows internally (see Section 6.2.6, “Access Control, Stage 1: Connection Verification”).

Suppose that a MySQL installation includes these two accounts:

```
-- create default proxy user
CREATE USER ''@''
  IDENTIFIED WITH some_plugin AS 'some_auth_string';
-- create anonymous user
CREATE USER ''@'%'
  IDENTIFIED BY 'some_password';
```

The first account (`''@''`) is intended as the default proxy user, used to authenticate connections for users who do not otherwise match a more-specific account. The second account (`''@'%'`) is an anonymous-user account, which might have been created, for example, to enable users without their own account to connect anonymously.

Both accounts have the same user part (`''`), which matches any user. And each account has a host part that matches any host. Nevertheless, there is a priority in account matching for connection attempts because the matching rules sort a host of `'%'` ahead of `''`. For accounts that do not match any more-specific account, the server attempts to authenticate them against `''@'%'` (the anonymous user) rather than `''@''` (the default proxy user). The result is that the default proxy account is never used.

To avoid this problem, use one of the following strategies:

- Remove the anonymous account so that it does not conflict with the default proxy user. This might be a good idea anyway if you want to associate every connection with a named user.
- Use a more-specific default proxy user that matches ahead of the anonymous user. For example, to permit only `localhost` proxy connections, use `''@'localhost'`:

  ```
  CREATE USER ''@'localhost'
    IDENTIFIED WITH some_plugin AS 'some_auth_string';
  ```

  In addition, modify any `GRANT PROXY` statements to name `''@'localhost'` rather than `''@''` as the proxy user.

Be aware that this strategy prevents anonymous-user connections from `localhost`.

- Create multiple proxy users, one for local connections and one for “everything else” (remote connections). This can be useful particularly when local users should have different privileges from remote users.

  Create the proxy users:

```
-- create proxy user for local connections
CREATE USER ''@'localhost'
  IDENTIFIED WITH some_plugin AS 'some_auth_string';
-- create proxy user for remote connections
CREATE USER ''@'%'
  IDENTIFIED WITH some_plugin AS 'some_auth_string';
```

  Create the proxied users:

```
-- create proxied user for local connections
CREATE USER 'developer'@'localhost'
  IDENTIFIED BY 'some_password';
-- create proxied user for remote connections
CREATE USER 'developer'@'%'
  IDENTIFIED BY 'some_password';
```

  Grant the proxy privilege to each proxy user for the corresponding proxied user:

```
GRANT PROXY ON 'developer'@'localhost' TO ''@'localhost';
GRANT PROXY ON 'developer'@'%' TO ''@'%';
```

  Finally, grant appropriate privileges to the local and remote proxied users (not shown).

  Assume that the `some_plugin`/`'some_auth_string'` combination causes `some_plugin` to map the client user name to `developer`. Local connections match the `''@'localhost'` proxy user, which maps to the `'developer'@'localhost'` proxied user. Remote connections match the `''@'%'` proxy user, which maps to the `'developer'@'%'` proxied user.

Server Support for Proxy User Mapping

Some authentication plugins implement proxy user mapping for themselves. For certain others, the MySQL server itself can map proxy users according to granted proxy privileges. If the `check_proxy_users` system variable is enabled, the server performs proxy user mapping for any authentication plugins that request it:

- By default, `check_proxy_users` is disabled, so the server performs no proxy user mapping even for authentication plugins that request it.
- With `check_proxy_users` enabled, it may also be necessary to enable plugin-specific system variables to take advantage of server proxy user mapping support:
  - For the `mysql_native_password` plugin, enable `mysql_native_password_proxy_users`.
  - For the `sha256_password` plugin, enable `sha256_password_proxy_users`.

Proxy user mapping performed by the server is subject to these restrictions:

- The server will not proxy to or from an anonymous user, even if the associated `PROXY` privilege is granted.
- When a single account has been granted proxy privileges for more than one proxied account, server proxy user mapping is nondeterministic. Therefore, granting to a single account proxy privileges for multiple proxied accounts is discouraged.

Proxy User System Variables

Two system variables help trace the proxy login process:

- `proxy_user`: This value is `NULL` if proxying is not used. Otherwise, it indicates the proxy user account. For example, if a client authenticates through the `''@''` proxy account, this variable is set as follows:

```
mysql> SELECT @@proxy_user;
+--------------+
| @@proxy_user |
+--------------+
| ''@''        |
+--------------+
```

- `external_user`: Sometimes the authentication plugin may use an external user to authenticate to the MySQL server. For example, when using Windows native authentication, a plugin that authenticates using the windows API does not need the login ID passed to it. However, it still uses a Windows user ID to authenticate. The plugin may return this external user ID (or the first 512 UTF-8 bytes of it) to the server using the `external_user` read-only session variable. If the plugin does not set this variable, its value is `NULL`.

6.3.12 User Account Locking

MySQL supports locking and unlocking user accounts using the `ACCOUNT LOCK` and `ACCOUNT UNLOCK` clauses for the `CREATE USER` and `ALTER USER` statements:

- When used with `CREATE USER`, these clauses specify the initial locking state for a new account. In the absence of either clause, the account is created in an unlocked state.
- When used with `ALTER USER`, these clauses specify the new locking state for an existing account. In the absence of either clause, the account locking state remains unchanged.

Account locking state is recorded in the `account_locked` column of the `mysql.user` table. The output from `SHOW CREATE USER` indicates whether an account is locked or unlocked.

If a client attempts to connect to a locked account, the attempt fails. The server increments the `Locked_connects` status variable that indicates the number of attempts to connect to a locked account, returns an `ER_ACCOUNT_HAS_BEEN_LOCKED` error, and writes a message to the error log:

```
Access denied for user 'user_name'@'host_name'.
Account is locked.
```

Locking an account does not affect being able to connect using a proxy user that assumes the identity of the locked account. It also does not affect the ability to execute stored programs or views that have a `DEFINER` clause naming the locked account. That is, the ability to use a proxied account or stored programs or views is not affected by locking the account.

The account-locking capability depends on the presence of the `account_locked` column in the `mysql.user` table. For upgrades to MySQL 5.7.6 and later from older versions, run `mysql_upgrade` to ensure that this column exists. For nonupgraded installations that have no `account_locked` column, the server treats all accounts as unlocked, and using the `ACCOUNT LOCK` or `ACCOUNT UNLOCK` clauses produces an error.

6.3.13 SQL-Based MySQL Account Activity Auditing

Applications can use the following guidelines to perform SQL-based auditing that ties database activity to MySQL accounts.

MySQL accounts correspond to rows in the `mysql.user` table. When a client connects successfully, the server authenticates the client to a particular row in this table. The `User` and `Host` column values

in this row uniquely identify the account and correspond to the `'user_name'@'host_name'` format in which account names are written in SQL statements.

The account used to authenticate a client determines which privileges the client has. Normally, the `CURRENT_USER()` function can be invoked to determine which account this is for the client user. Its value is constructed from the `User` and `Host` columns of the `user` table row for the account.

However, there are circumstances under which the `CURRENT_USER()` value corresponds not to the client user but to a different account. This occurs in contexts when privilege checking is not based the client's account:

- Stored routines (procedures and functions) defined with the `SQL SECURITY DEFINER` characteristic
- Views defined with the `SQL SECURITY DEFINER` characteristic
- Triggers and events

In those contexts, privilege checking is done against the `DEFINER` account and `CURRENT_USER()` refers to that account, not to the account for the client who invoked the stored routine or view or who caused the trigger to activate. To determine the invoking user, you can call the `USER()` function, which returns a value indicating the actual user name provided by the client and the host from which the client connected. However, this value does not necessarily correspond directly to an account in the `user` table, because the `USER()` value never contains wildcards, whereas account values (as returned by `CURRENT_USER()`) may contain user name and host name wildcards.

For example, a blank user name matches any user, so an account of `''@'localhost'` enables clients to connect as an anonymous user from the local host with any user name. In this case, if a client connects as `user1` from the local host, `USER()` and `CURRENT_USER()` return different values:

```
mysql> SELECT USER(), CURRENT_USER();
+----------------+----------------+
| USER()         | CURRENT_USER() |
+----------------+----------------+
| user1@localhost | @localhost     |
+----------------+----------------+
```

The host name part of an account can contain wildcards, too. If the host name contains a `'%'` or `'_'` pattern character or uses netmask notation, the account can be used for clients connecting from multiple hosts and the `CURRENT_USER()` value will not indicate which one. For example, the account `'user2'@'%.example.com'` can be used by `user2` to connect from any host in the `example.com` domain. If `user2` connects from `remote.example.com`, `USER()` and `CURRENT_USER()` return different values:

```
mysql> SELECT USER(), CURRENT_USER();
+--------------------------+--------------------+
| USER()                   | CURRENT_USER()     |
+--------------------------+--------------------+
| user2@remote.example.com | user2@%.example.com |
+--------------------------+--------------------+
```

If an application must invoke `USER()` for user auditing (for example, if it does auditing from within triggers) but must also be able to associate the `USER()` value with an account in the `user` table, it is necessary to avoid accounts that contain wildcards in the `User` or `Host` column. Specifically, do not permit `User` to be empty (which creates an anonymous-user account), and do not permit pattern characters or netmask notation in `Host` values. All accounts must have a nonempty `User` value and literal `Host` value.

With respect to the previous examples, the `''@'localhost'` and `'user2'@'%.example.com'` accounts should be changed not to use wildcards:

```
RENAME USER ''@'localhost' TO 'user1'@'localhost';
RENAME USER 'user2'@'%.example.com' TO 'user2'@'remote.example.com';
```

If `user2` must be able to connect from several hosts in the `example.com` domain, there should be a separate account for each host.

To extract the user name or host name part from a `CURRENT_USER()` or `USER()` value, use the `SUBSTRING_INDEX()` function:

```
mysql> SELECT SUBSTRING_INDEX(CURRENT_USER(),'@',1);
+--------------------------------------+
| SUBSTRING_INDEX(CURRENT_USER(),'@',1) |
+--------------------------------------+
| user1                                |
+--------------------------------------+

mysql> SELECT SUBSTRING_INDEX(CURRENT_USER(),'@',-1);
+---------------------------------------+
| SUBSTRING_INDEX(CURRENT_USER(),'@',-1) |
+---------------------------------------+
| localhost                             |
+---------------------------------------+
```

6.4 Using Secure Connections

With an unencrypted connection between the MySQL client and the server, someone with access to the network could watch all your traffic and inspect the data being sent or received between client and server.

When you must move information over a network in a secure fashion, an unencrypted connection is unacceptable. To make any kind of data unreadable, use encryption. Encryption algorithms must include security elements to resist many kinds of known attacks such as changing the order of encrypted messages or replaying data twice.

MySQL supports secure (encrypted) connections between clients and the server using the TLS (Transport Layer Security) protocol. TLS is sometimes referred to as SSL (Secure Sockets Layer) but MySQL does not actually use the SSL protocol for secure connections because it provides weak encryption (see Section 6.4.3, “Secure Connection Protocols and Ciphers”).

TLS uses encryption algorithms to ensure that data received over a public network can be trusted. It has mechanisms to detect data change, loss, or replay. TLS also incorporates algorithms that provide identity verification using the X509 standard.

X509 makes it possible to identify someone on the Internet. In basic terms, there should be some entity called a “Certificate Authority” (or CA) that assigns electronic certificates to anyone who needs them. Certificates rely on asymmetric encryption algorithms that have two encryption keys (a public key and a secret key). A certificate owner can present the certificate to another party as proof of identity. A certificate consists of its owner's public key. Any data encrypted using this public key can be decrypted only using the corresponding secret key, which is held by the owner of the certificate.

MySQL can be compiled for secure-connection support using OpenSSL or yaSSL. For a comparison of the two packages, see Section 6.4.1, “OpenSSL Versus yaSSL” For information about the encryption protocols and ciphers each package supports, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

MySQL performs encryption on a per-connection basis, and use of encryption for a given user can be optional or mandatory. This enables you to choose an encrypted or unencrypted connection according to the requirements of individual applications. For information on how to require users to use encrypted connections, see the discussion of the `REQUIRE` clause of the `CREATE USER` statement in Section 13.7.1.3, “CREATE USER Syntax”. See also the description of the `require_secure_transport` system variable at Section 5.1.5, “Server System Variables”

Secure connections are available through the MySQL C API using the `mysql_ssl_set()` and `mysql_options()` functions. See Section 27.7.7.75, “mysql_ssl_set()”, and Section 27.7.7.50, “mysql_options()”.

Replication uses the C API, so secure connections can be used between master and slave servers. See Section 18.3.9, “Setting Up Replication to Use Secure Connections”.

It is also possible to connect securely from within an SSH connection to the MySQL server host. For an example, see Section 6.4.7, “Connecting to MySQL Remotely from Windows with SSH”.

6.4.1 OpenSSL Versus yaSSL

MySQL can be compiled using OpenSSL or yaSSL, both of which enable secure conections based on the OpenSSL API:

- MySQL Enterprise Edition binary distributions are compiled using OpenSSL. It is not possible to use yaSSL with MySQL Enterprise Edition.
- MySQL Community Edition binary distributions are compiled using yaSSL.
- MySQL Community Edition source distributions can be compiled using either OpenSSL or yaSSL (see Section 6.4.2, “Building MySQL with Support for Secure Connections”).

OpenSSL and yaSSL offer the same basic functionality, but MySQL distributions compiled using OpenSSL have additional features:

- OpenSSL supports a wider range of encryption ciphers from which to choose for the `--ssl-cipher` option. OpenSSL supports the `--ssl-capath`, `--ssl-crl`, and `--ssl-crlpath` options. See Section 6.4.5, “Command Options for Secure Connections”.
- Accounts that authenticate using the `sha256_password` plugin can use RSA key files for secure password exchange over unencrypted connections. See Section 6.5.1.2, “SHA-256 Pluggable Authentication”.
- The server can automatically generate missing SSL and RSA certificate and key files at startup. See Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.
- OpenSSL supports more encryption modes for the `AES_ENCRYPT()` and `AES_DECRYPT()` functions. See Section 12.13, “Encryption and Compression Functions”

Certain OpenSSL-related system and status variables are present only if MySQL was compiled using OpenSSL:

- `auto_generate_certs`
- `sha256_password_auto_generate_rsa_keys`
- `sha256_password_private_key_path`
- `sha256_password_public_key_path`
- `Rsa_public_key`

To determine whether your server was compiled using OpenSSL, test the existence of any of those variables. For example, this statement returns a row if OpenSSL was used and an empty result if yaSSL was used:

```
SHOW STATUS LIKE 'Rsa_public_key';
```

6.4.2 Building MySQL with Support for Secure Connections

To use SSL connections between the MySQL server and client programs, your system must support either OpenSSL or yaSSL:

- MySQL Enterprise Edition binary distributions are compiled using OpenSSL. It is not possible to use yaSSL with MySQL Enterprise Edition.
- MySQL Community Edition binary distributions are compiled using yaSSL.
- MySQL Community Edition source distributions can be compiled using either OpenSSL or yaSSL.

If you compile MySQL from a source distribution, `CMake` configures the distribution to use yaSSL by default. To compile using OpenSSL instead, use this procedure:

1. Ensure OpenSSL 1.0.1 or higher is installed on your system. To obtain OpenSSL, visit http://www.openssl.org.

   If the installed OpenSSL version is lower than 1.0.1, `CMake` produces an error at MySQL configuration time.

2. To use OpenSSL, add the `-DWITH_SSL=system` option to the `CMake` command you normally use to configure the MySQL source distribution. For example:

   ```
   shell> cmake . -DWITH_SSL=system
   ```

   That command configures the distribution to use the installed OpenSSL library. Alternatively, to explicitly specify the path name to the OpenSSL installation, use the following syntax. This can be useful if you have multiple versions of OpenSSL installed, to prevent `CMake` from choosing the wrong one:

   ```
   shell> cmake . -DWITH_SSL=path_name
   ```

   See Section 2.8.4, “MySQL Source-Configuration Options”.

3. Compile and install the distribution.

To check whether a `mysqld` server supports secure connections, examine the value of the `have_ssl` system variable:

```
mysql> SHOW VARIABLES LIKE 'have_ssl';
+---------------+-------+
| Variable_name | Value |
+---------------+-------+
| have_ssl      | YES   |
+---------------+-------+
```

If the value is `YES`, the server supports secure connections. If the value is `DISABLED`, the server is capable of supporting secure connections but was not started with the appropriate `--ssl-xxx` options to enable secure connections to be used; see Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

To determine whether a server was compiled using OpenSSL or yaSSL, check the existence of any of the system or status variables that are present only for OpenSSL. See Section 6.4.1, “OpenSSL Versus yaSSL”

6.4.3 Secure Connection Protocols and Ciphers

To determine which encryption protocol and cipher are in use for an encrypted connection, use the following statements to check the values of the `Ssl_version` and `Ssl_cipher` status variables:

```
mysql> SHOW SESSION STATUS LIKE 'Ssl_version';
+---------------+-------+
| Variable_name | Value |
+---------------+-------+
| Ssl_version   | TLSv1 |
+---------------+-------+
mysql> SHOW SESSION STATUS LIKE 'Ssl_cipher';
+---------------+--------------------------+
| Variable_name | Value                    |
+---------------+--------------------------+
| Ssl_cipher    | DHE-RSA-AES128-GCM-SHA256 |
+---------------+--------------------------+
```

If the connection is not encrypted, both variables have an empty value.

MySQL supports encrypted connections using TLS protocols:

- When compiled using OpenSSL 1.0.1 or higher, MySQL supports the TLSv1, TLSv1.1, and TLSv1.2 protocols.
- When compiled using the bundled version of yaSSL, MySQL supports the TLSv1 and TLSv1.1 protocols.

The value of the `tls_version` system variable determines which protocols the server is permitted to use from those that are available. The `tls_version` value is a comma-separated list containing one or more of these protocols (not case sensitive): TLSv1, TLSv1.1, TLSv1.2. By default, this variable lists all protocols supported by the SSL library used to compile MySQL (`TLSv1,TLSv1.1,TLSv1.2` for OpenSSL, `TLSv1,TLSv1.1` for yaSSL). To determine the value of `tls_version` at runtime, use this statement:

```
mysql> SHOW GLOBAL VARIABLES LIKE 'tls_version';
+---------------+-----------------------+
| Variable_name | Value                 |
+---------------+-----------------------+
| tls_version   | TLSv1,TLSv1.1,TLSv1.2 |
+---------------+-----------------------+
```

To change the value of `tls_version`, set it at server startup. For example, to prohibit connections that use the less-secure TLSv1 protocol, use these lines in the server `my.cnf` file:

```
[mysqld]
tls_version=TLSv1.1,TLSv1.2
```

To be even more restrict and permit only TLSv1.2 connections, set `tls_version` like this (assuming that your server is compiled using OpenSSL because yaSSL does not support TLSv1.2):

```
[mysqld]
tls_version=TLSv1.2
```

For client programs, the `--tls-version` option enables specifying the TLS protocols permitted per client invocation. The value format is the same as for `tls_version`.

By default, MySQL attempts to use the highest TLS protocol version available, depending on which SSL library was used to compile the server and client, which key size is used, and whether the server or client are restricted from using some protocols; for example, by means of `tls_version`/`--tls-version`:

- If the server and client are compiled using OpenSSL, TLSv1.2 is used if possible.
- If either or both the server and client are compiled using yaSSL, TLSv1.1 is used if possible.

- TLSv1.2 does not work with all ciphers that have a key size of 512 bits or less. To use this protocol with such a key, use `--ssl-cipher` to specify the cipher name explicitly:

```
AES128-SHA
AES128-SHA256
AES256-SHA
AES256-SHA256
CAMELLIA128-SHA
CAMELLIA256-SHA
DES-CBC3-SHA
DHE-RSA-AES256-SHA
RC4-MD5
RC4-SHA
SEED-SHA
```

- For better security, use a certificate with an RSA key size of of 2048 bits or more.

If the server and client protocol capabilities have no protocol in common, the server terminates the connection request. For example, if the server is configured with `tls_version=TLSv1.1,TLSv1.2`, connection attempts will fail for clients invoked with `--tls-version=TLSv1`, and for older clients that do not support the `--tls-version` option and implicitly support only TLSv1.

To determine which ciphers a given server supports, use the following statement to check the value of the `Ssl_cipher_list` status variable:

```
SHOW SESSION STATUS LIKE 'Ssl_cipher_list';
```

The set of available ciphers depends on your MySQL version and whether MySQL was compiled using OpenSSL or yaSSL, and (for OpenSSL) the library version used to compile MySQL.

Order of ciphers passed by MySQL to the SSL library is significant. More secure ciphers are mentioned first in the list, and the first cipher supported by the provided certificate is selected.

MySQL passes this cipher list to OpenSSL:

```
ECDHE-ECDSA-AES128-GCM-SHA256
ECDHE-ECDSA-AES256-GCM-SHA384
ECDHE-RSA-AES128-GCM-SHA256
ECDHE-RSA-AES256-GCM-SHA384
ECDHE-ECDSA-AES128-SHA256
ECDHE-RSA-AES128-SHA256
ECDHE-ECDSA-AES256-SHA384
ECDHE-RSA-AES256-SHA384
DHE-RSA-AES128-GCM-SHA256
DHE-DSS-AES128-GCM-SHA256
DHE-RSA-AES128-SHA256
DHE-DSS-AES128-SHA256
DHE-DSS-AES256-GCM-SHA384
DHE-RSA-AES256-SHA256
DHE-DSS-AES256-SHA256
ECDHE-RSA-AES128-SHA
ECDHE-ECDSA-AES128-SHA
ECDHE-RSA-AES256-SHA
ECDHE-ECDSA-AES256-SHA
DHE-DSS-AES128-SHA
DHE-RSA-AES128-SHA
TLS_DHE_DSS_WITH_AES_256_CBC_SHA
DHE-RSA-AES256-SHA
AES128-GCM-SHA256
DH-DSS-AES128-GCM-SHA256
ECDH-ECDSA-AES128-GCM-SHA256
AES256-GCM-SHA384
DH-DSS-AES256-GCM-SHA384
ECDH-ECDSA-AES256-GCM-SHA384
AES128-SHA256
DH-DSS-AES128-SHA256
```

```
ECDH-ECDSA-AES128-SHA256
AES256-SHA256
DH-DSS-AES256-SHA256
ECDH-ECDSA-AES256-SHA384
AES128-SHA
DH-DSS-AES128-SHA
ECDH-ECDSA-AES128-SHA
AES256-SHA
DH-DSS-AES256-SHA
ECDH-ECDSA-AES256-SHA
DHE-RSA-AES256-GCM-SHA384
DH-RSA-AES128-GCM-SHA256
ECDH-RSA-AES128-GCM-SHA256
DH-RSA-AES256-GCM-SHA384
ECDH-RSA-AES256-GCM-SHA384
DH-RSA-AES128-SHA256
ECDH-RSA-AES128-SHA256
DH-RSA-AES256-SHA256
ECDH-RSA-AES256-SHA384
ECDHE-RSA-AES128-SHA
ECDHE-ECDSA-AES128-SHA
ECDHE-RSA-AES256-SHA
ECDHE-ECDSA-AES256-SHA
DHE-DSS-AES128-SHA
DHE-RSA-AES128-SHA
TLS_DHE_DSS_WITH_AES_256_CBC_SHA
DHE-RSA-AES256-SHA
AES128-SHA
DH-DSS-AES128-SHA
ECDH-ECDSA-AES128-SHA
AES256-SHA
DH-DSS-AES256-SHA
ECDH-ECDSA-AES256-SHA
DH-RSA-AES128-SHA
ECDH-RSA-AES128-SHA
DH-RSA-AES256-SHA
ECDH-RSA-AES256-SHA
DES-CBC3-SHA
```

MySQL passes this cipher list to yaSSL:

```

DHE-RSA-AES256-SHA
DHE-RSA-AES128-SHA
AES128-RMD
DES-CBC3-RMD
DHE-RSA-AES256-RMD
DHE-RSA-AES128-RMD
DHE-RSA-DES-CBC3-RMD
AES256-SHA
RC4-SHA
RC4-MD5
DES-CBC3-SHA
DES-CBC-SHA
EDH-RSA-DES-CBC3-SHA
EDH-RSA-DES-CBC-SHA
AES128-SHA:AES256-RMD
```

These cipher restrictions are in place:

- The following ciphers are permanently restricted:

  ```

  !DHE-DSS-DES-CBC3-SHA
  !DHE-RSA-DES-CBC3-SHA
  !ECDH-RSA-DES-CBC3-SHA
  !ECDH-ECDSA-DES-CBC3-SHA
  !ECDHE-RSA-DES-CBC3-SHA
  !ECDHE-ECDSA-DES-CBC3-SHA
  ```

- The following categories of ciphers are permanently restricted:

```
!aNULL
!eNULL
!EXPORT
!LOW
!MD5
!DES
!RC2
!RC4
!PSK
!SSLv3
```

If the server is started using a compatible certificate that uses any of the preceding restricted ciphers or cipher categories, the server starts with support for secure connections disabled.

6.4.4 Configuring MySQL to Use Secure Connections

To enable secure connections, the proper options must be used to specify the appropriate certificate and key files. For a complete list of options related to establishment of secure connections, see Section 6.4.5, “Command Options for Secure Connections”.

If you need to create the required certificate and key files, see Section 6.4.6, “Creating SSL and RSA Certificates and Keys”.

- Server-Side Configuration for Secure Connections
- Client-Side Configuration for Secure Connections
- C API Configuration for Secure Connections

Server-Side Configuration for Secure Connections

To start the MySQL server so that it permits clients to connect securely, use options that identify the certificate and key files the server uses when establishing a secure connection:

- `--ssl-ca` identifies the Certificate Authority (CA) certificate.
- `--ssl-cert` identifies the server public key certificate. This can be sent to the client and authenticated against the CA certificate that it has.
- `--ssl-key` identifies the server private key.

For example, start the server with these lines in the `my.cnf` file, changing the file names as necessary:

```
[mysqld]
ssl-ca=ca.pem
ssl-cert=server-cert.pem
ssl-key=server-key.pem
```

Each option names a file in PEM format. If you have a MySQL source distribution, you can test your setup using the demonstration certificate and key files in its `mysql-test/std_data` directory.

The server-side `--ssl` option value is enabled by default.

MySQL servers compiled using OpenSSL can generate missing certificate and key files automatically at startup. See Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.

The server performs certificate and key file autodiscovery. If `--ssl` is enabled (possibly along with `--ssl-cipher`) and other `--ssl-xxx` options are not given to configure secure connections explicitly, the server attempts to enable support for secure connections automatically at startup:

- If the server discovers valid certificate and key files named `ca.pem`, `server-cert.pem`, and `server-key.pem` in the data directory, it enables support for secure connections by clients. (The files need not have been autogenerated; what matters is that they have the indicated names and are valid.)
- If the server does not find valid certificate and key files in the data directory, it continues executing but does not enable secure connections.

If the server automatically enables support for secure connections, it writes a message to the error log. If the server discovers that the CA certificate is self-signed, it writes a warning to the error log. (The certificate will be self-signed if created automatically by the server or manually using `mysql_ssl_rsa_setup`.)

For any certificate and key files that the server discovers and uses automatically, it uses the file names to set the corresponding system variables (`ssl_ca`, `ssl_cert`, `ssl_key`).

For further control over whether clients must connect securely, use the `require_secure_transport` system variable; see Section 5.1.5, “Server System Variables”. For information about permitted encryption protocols and ciphers, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

Client-Side Configuration for Secure Connections

For client programs, options for secure connections are similar to those used on the server side, but `--ssl-cert` and `--ssl-key` identify the client public and private key:

- `--ssl-ca` identifies the Certificate Authority (CA) certificate. This option, if used, must specify the same certificate used by the server.
- `--ssl-cert` identifies the client public key certificate.
- `--ssl-key` identifies the client private key.

To connect securely to a MySQL server that supports secure connections, the options that a client must specify depend on the encryption requirements of the MySQL account used by the client. (See the discussion of the `REQUIRE` clause in Section 13.7.1.3, “CREATE USER Syntax”.)

Suppose that you want to connect using an account that has no special encryption requirements or was created using a `CREATE USER` statement that includes the `REQUIRE SSL` option. As a recommended set of secure-connection options, start the server with at least `--ssl-cert` and `--ssl-key`, and invoke the client with `--ssl-ca`. A client can connect securely like this:

```
shell> mysql --ssl-ca=ca.pem
```

To require that a client certificate also be specified, create the account using the `REQUIRE X509` option. Then the client must also specify the proper client key and certificate files or the server will reject the connection:

```
shell> mysql --ssl-ca=ca.pem \
       --ssl-cert=client-cert.pem \
       --ssl-key=client-key.pem
```

To prevent use of encryption and override other `--ssl-xxx` options, invoke the client program with `--ssl-mode=DISABLED`.

```
shell> mysql --ssl-mode=DISABLED
```

Client programs attempt to establish a secure connection by default whenever the server supports secure connections:

- In the absence of an `--ssl-mode` or `--ssl` option, the client falls back to an unencrypted connection if a secure connection cannot be established.
- To require a secure connection and fail if one cannot be established, invoke the client with `--ssl-mode=REQUIRED`.
- To use an unencrypted connection, invoke the client with `--ssl-mode=DISABLED`.

For information about permitted encryption protocols and ciphers, see Section 6.4.3, "Secure Connection Protocols and Ciphers".

A client can determine whether the current connection with the server uses encryption by checking the value of the `Ssl_cipher` status variable. If the value is empty, the connection is not encrypted. Otherwise, the connection is encrypted and the value indicates the encryption cipher. For example:

```
mysql> SHOW STATUS LIKE 'Ssl_cipher';
+---------------+--------------------+
| Variable_name | Value              |
+---------------+--------------------+
| Ssl_cipher    | DHE-RSA-AES256-SHA |
+---------------+--------------------+
```

For the `mysql` client, an alternative is to use the `STATUS` or `\s` command and check the `SSL` line:

```
mysql> \s
...
SSL: Cipher in use is DHE-RSA-AES256-SHA
...
```

Or:

```
mysql> \s
...
SSL: Not in use
...
```

C API Configuration for Secure Connections

The C API enables application programs to use secure connections:

- To establish a secure connection, use the `mysql_ssl_set()` C API function to set the appropriate certificate options before calling `mysql_real_connect()`. See Section 27.7.7.75, "mysql_ssl_set()". To require the use of a secure connection, call `mysql_options()` with the `MYSQL_OPT_SSL_MODE` option. To establish permitted encryption protocols, call `mysql_options()` with the `MYSQL_OPT_TLS_VERSION` option.
- To determine whether encryption is in use after the connection is established, use `mysql_get_ssl_cipher()`. A non-`NULL` return value indicates an encrypted connection and names the cipher used for encryption. A `NULL` return value indicates that encryption is not being used. See Section 27.7.7.34, "mysql_get_ssl_cipher()".

Replication uses the C API, so secure connections can be used between master and slave servers. See Section 18.3.9, "Setting Up Replication to Use Secure Connections".

6.4.5 Command Options for Secure Connections

This section describes options that specify whether to use secure connections and the names of certificate and key files. These options can be given on the command line or in an option file. For examples of suggested use and how to check whether a connection is secure, see Section 6.4.4, "Configuring MySQL to Use Secure Connections".

Table 6.9 Secure-Connection Option Summary

Format	Description
--skip-ssl	Do not use secure connection
--ssl	Enable secure connection
--ssl-ca	Path of file that contains list of trusted SSL CAs
--ssl-capath	Path of directory that contains trusted SSL CA certificates in PEM format
--ssl-cert	Path of file that contains X509 certificate in PEM format
--ssl-cipher	List of permitted ciphers to use for connection encryption
--ssl-crl	Path of file that contains certificate revocation lists
--ssl-crlpath	Path of directory that contains certificate revocation list files
--ssl-key	Path of file that contains X509 key in PEM format
--ssl-mode	Security state of connection to server
--tls-version	Protocols permitted for secure connections

- `--ssl`

Note

The client-side `--ssl` option is removed in MySQL 8.0. For client programs, use `--ssl-mode` instead.

For the MySQL server, this option specifies that the server permits but does not require secure connections. The option is enabled on the server side by default.

MySQL servers compiled using OpenSSL can generate missing certificate and key files automatically at startup. See Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.

The server performs certificate and key file autodiscovery. If `--ssl` is enabled (possibly along with `--ssl-cipher`) and other `--ssl-xxx` options are not given to configure secure connections explicitly, the server attempts to enable support for secure connections automatically at startup:

- If the server discovers valid certificate and key files named `ca.pem`, `server-cert.pem`, and `server-key.pem` in the data directory, it enables support for secure connections by clients. (The files need not have been autogenerated; what matters is that they have the indicated names and are valid.)
- If the server does not find valid certificate and key files in the data directory, it continues executing but does not enable secure connections.

As a recommended set of options to enable secure connections, use at least `--ssl-cert` and `--ssl-key` on the server side and `--ssl-ca` on the client side. See Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

`--ssl` is implied by other `--ssl-xxx` options, as indicated in the descriptions for those options.

The `--ssl` option in negated form overrides other `--ssl-xxx` options and indicates that encryption should *not* be used. To do this, specify the option as `--ssl=0` or a synonym (`--skip-ssl`, `--disable-ssl`).

To require use of secure connections by a MySQL account, use `CREATE USER` to create the account with at least a `REQUIRE SSL` clause, or use `ALTER USER` for an existing account to add a `REQUIRE` clause. Connections for the account will be rejected unless MySQL supports secure connections and the server and client have been started with the proper secure-connection options.

The `REQUIRE` clause permits other encryption-related options, which can be used to enforce stricter requirements than `REQUIRE SSL`. For additional details about which command options may or must be specified by clients that connect using accounts configured using the various `REQUIRE` options, see the description of `REQUIRE` in Section 13.7.1.3, “CREATE USER Syntax”.

- `--ssl-ca=file_name`

  The path to a file in PEM format that contains a list of trusted SSL certificate authorities. On the server side, this option implies `--ssl`.

  If you use encryption when establishing a client connection, to tell the client not to authenticate the server certificate, specify neither `--ssl-ca` nor `--ssl-capath`. The server still verifies the client according to any applicable requirements established for the client account, and it still uses any `--ssl-ca` or `--ssl-capath` option values specified at server startup.

- `--ssl-capath=dir_name`

  The path to a directory that contains trusted SSL certificate authority certificates in PEM format. On the server side, this option implies `--ssl`.

  If you use encryption when establishing a client connection, to tell the client not to authenticate the server certificate, specify neither `--ssl-ca` nor `--ssl-capath`. The server still verifies the client according to any applicable requirements established for the client account, and it still uses any `--ssl-ca` or `--ssl-capath` option values specified at server startup.

  MySQL distributions compiled using OpenSSL support the `--ssl-capath` option (see Section 6.4.1, “OpenSSL Versus yaSSL”). Distributions compiled using yaSSL do not because yaSSL does not look in any directory and does not follow a chained certificate tree. yaSSL requires that all components of the CA certificate tree be contained within a single CA certificate tree and that each certificate in the file has a unique SubjectName value. To work around this yaSSL limitation, concatenate the individual certificate files comprising the certificate tree into a new file and specify that file as the value of the `--ssl-ca` option.

- `--ssl-cert=file_name`

  The name of the SSL certificate file in PEM format to use for establishing a secure connection. On the server side, this option implies `--ssl`.

- `--ssl-cipher=cipher_list`

  A list of permissible ciphers to use for connection encryption. If no cipher in the list is supported, encrypted connections will not work. On the server side, this option implies `--ssl`.

  For greatest portability, *`cipher_list`* should be a list of one or more cipher names, separated by colons. This format is understood both by OpenSSL and yaSSL. Examples:

  ```
  --ssl-cipher=AES128-SHA
  --ssl-cipher=DHE-RSA-AES256-SHA:AES128-SHA
  ```

  OpenSSL supports a more flexible syntax for specifying ciphers, as described in the OpenSSL documentation at https://www.openssl.org/docs/manmaster/man1/ciphers.html. yaSSL does not, so attempts to use that extended syntax fail for a MySQL distribution compiled using yaSSL.

  For information about which encryption ciphers MySQL supports, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `--ssl-crl=file_name`

  The path to a file containing certificate revocation lists in PEM format. On the server side, this option implies `--ssl`.

If neither `--ssl-crl` nor `--ssl-crlpath` is given, no CRL checks are performed, even if the CA path contains certificate revocation lists.

MySQL distributions compiled using OpenSSL support the `--ssl-crl` option (see Section 6.4.1, “OpenSSL Versus yaSSL”). Distributions compiled using yaSSL do not because revocation lists do not work with yaSSL.

- `--ssl-crlpath=`*`dir_name`*

  The path to a directory that contains files containing certificate revocation lists in PEM format. On the server side, this option implies `--ssl`.

  If neither `--ssl-crl` nor `--ssl-crlpath` is given, no CRL checks are performed, even if the CA path contains certificate revocation lists.

  MySQL distributions compiled using OpenSSL support the `--ssl-crlpath` option (see Section 6.4.1, “OpenSSL Versus yaSSL”). Distributions compiled using yaSSL do not because revocation lists do not work with yaSSL.

- `--ssl-key=`*`file_name`*

  The name of the SSL key file in PEM format to use for establishing a secure connection. On the server side, this option implies `--ssl`.

  If the key file is protected by a passphrase, the program prompts the user for the passphrase. The password must be given interactively; it cannot be stored in a file. If the passphrase is incorrect, the program continues as if it could not read the key.

  For better security, use a certificate with an RSA key size of of 2048 bits or more.

- `--ssl-mode=`*`mode`*

  This option is available only for client programs, not the server. It specifies the security state of the connection to the server. The following option values are permitted:

  - `PREFERRED`: Establish a secure (encrypted) connection if the server supports secure connections. Fall back to an unencrypted connection otherwise. This is the default if `--ssl-mode` is not specified.
  - `DISABLED`: Establish an unencrypted connection.
  - `REQUIRED`: Establish a secure connection if the server supports secure connections. The connection attempt fails if a secure connection cannot be established.
  - `VERIFY_CA`: Like `REQUIRED`, but additionally verify the server TLS certificate against the configured Certificate Authority (CA) certificates. The connection attempt fails if no valid matching CA certificates are found.
  - `VERIFY_IDENTITY`: Like `VERIFY_CA`, but additionally verify that the server certificate matches the host to which the connection is attempted.

  Use of the `--ssl-ca` or `--ssl-capath` option implies `--ssl-mode=VERIFY_CA`, if `--ssl-mode` is not explicitly set otherwise.

  If `--ssl-mode` is explicit, use of a value other than `VERIFY_CA` or `VERIFY_IDENTITY` with an explicit `--ssl-ca` or `--ssl-capath` option produces a warning that no verification of the server certificate will be done, despite CA certificate options being specified.

  To require use of secure connections by a MySQL account, use `CREATE USER` to create the account with at least a `REQUIRE SSL` clause, or use `ALTER USER` for an existing account to add a `REQUIRE`

clause. Connections for the account will be rejected unless MySQL supports secure connections and the server and client have been started with the proper secure-connection options.

The `REQUIRE` clause permits other encryption-related options, which can be used to enforce stricter requirements than `REQUIRE SSL`. For additional details about which command options may or must be specified by clients that connect using accounts configured using the various `REQUIRE` options, see the description of `REQUIRE` in Section 13.7.1.3, “CREATE USER Syntax”.

- `--tls-version=`*`protocol_list`*

  For client programs, the protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

  On the server side, the `tls_version` system variable can be used instead.

6.4.6 Creating SSL and RSA Certificates and Keys

The following discussion describes how to create the files required for SSL and RSA support in MySQL. File creation can be performed using facilities provided by MySQL itself, or by invoking the `openssl` command directly.

SSL certificate and key files enable MySQL to support secure connections using SSL. See Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

RSA key files enable MySQL to support secure password exchange over unencrypted connections for accounts authenticated by the `sha256_password` plugin. See Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

6.4.6.1 Creating SSL and RSA Certificates and Keys using MySQL

MySQL provides these ways to create the SSL certificate and key files and RSA key-pair files required to support secure connections using SSL and secure password exchange using RSA over unencrypted connections, if those files are missing:

- The server can autogenerate these files at startup, for MySQL distributions compiled using OpenSSL.
- Users can invoke the `mysql_ssl_rsa_setup` utility manually.
- For some distribution types, such as RPM packages, `mysql_ssl_rsa_setup` invocation occurs during data directory initialization. In this case, the MySQL distribution need not have been compiled using OpenSSL as long as the `openssl` command is available.

Important

Server autogeneration and `mysql_ssl_rsa_setup` help lower the barrier to using SSL by making it easier to generate the required files. However, certificates generated by these methods are self-signed, which may not be very secure. After you gain experience using such files, consider obtaining certificate/key material from a registered certificate authority.

- Automatic Generation of SSL and RSA Files
- Manual Generation of SSL and RSA Files Using mysql_ssl_rsa_setup
- SSL and RSA File Characteristics

Automatic Generation of SSL and RSA Files

MySQL servers have the capability of automatically generating missing SSL and RSA files at startup, for MySQL distributions compiled using OpenSSL. The `auto_generate_certs` and

`sha256_password_auto_generate_rsa_keys` system variables control automatic generation of these files. Both variables are enabled by default. They can be enabled at startup and inspected but not set at runtime.

At startup, the server automatically generates server-side and client-side SSL certificate and key files in the data directory if the `auto_generate_certs` system variable is enabled, no SSL options other than `--ssl` are specified, and the server-side SSL files are missing from the data directory. These files enable secure client connections using SSL; see Section 6.4.4, "Configuring MySQL to Use Secure Connections".

1. The server checks the data directory for SSL files with the following names:

```
ca.pem
server-cert.pem
server-key.pem
```

2. If any of those files are present, the server creates no SSL files. Otherwise, it creates them, plus some additional files:

```
ca.pem               Self-signed CA certificate
ca-key.pem           CA private key
server-cert.pem      Server certificate
server-key.pem       Server private key
client-cert.pem      Client certificate
client-key.pem       Client private key
```

3. If the server autogenerates SSL files, it uses the names of the `ca.pem`, `server-cert.pem`, and `server-key.pem` files to set the corresponding system variables (`ssl_ca`, `ssl_cert`, `ssl_key`).

At startup, the server automatically generates RSA private/public key-pair files in the data directory if the `sha256_password_auto_generate_rsa_keys` system variable is enabled, no RSA options are specified, and the RSA files are missing from the data directory. These files enable secure password exchange using RSA over unencrypted connections for accounts authenticated by the `sha256_password` plugin; see Section 6.5.1.2, "SHA-256 Pluggable Authentication".

1. The server checks the data directory for RSA files with the following names:

```
private_key.pem      Private member of private/public key pair
public_key.pem       Public member of private/public key pair
```

2. If any of these files are present, the server creates no RSA files. Otherwise, it creates them.

3. If the server autogenerates the RSA files, it uses their names to set the corresponding system variables (`sha256_password_private_key_path`, `sha256_password_public_key_path`).

Manual Generation of SSL and RSA Files Using mysql_ssl_rsa_setup

MySQL distributions include a `mysql_ssl_rsa_setup` utility that can be invoked manually to generate SSL and RSA files. This utility is included with all MySQL distributions (whether compiled using OpenSSL or yaSSL), but it does require that the `openssl` command be available. For usage instructions, see Section 4.4.3, "`mysql_ssl_rsa_setup` — Create SSL/RSA Files".

SSL and RSA File Characteristics

SSL and RSA files created automatically by the server or by invoking `mysql_ssl_rsa_setup` have these characteristics:

- SSL and RSA keys are 2048 bit.
- The SSL CA certificate is self signed.
- The SSL server and client certificates are signed with the CA certificate and key, using the `sha256WithRSAEncryption` signature algorithm.

- SSL certificates use these Common Name (CN) values, with the appropriate certificate type (CA, Server, Client):

```
ca.pem:         MySQL_Server_suffix_Auto_Generated_CA_Certificate
server-cert.pm: MySQL_Server_suffix_Auto_Generated_Server_Certificate
client-cert.pm: MySQL_Server_suffix_Auto_Generated_Client_Certificate
```

  The *`suffix`* value is based on the MySQL version number. For files generated by `mysql_ssl_rsa_setup`, the suffix can be specified explicitly using the `--suffix` option.

  For files generated by the server, if the resulting CN values exceed 64 characters, the *`_suffix`* portion of the name is omitted.

- SSL files have blank values for Country (C), State or Province (ST), Organization (O), Organization Unit Name (OU) and email address.
- SSL files created by the server or by `mysql_ssl_rsa_setup` are valid for ten years from the time of generation.
- RSA files do not expire.
- SSL files have different serial numbers for each certificate/key pair (1 for CA, 2 for Server, 3 for Client).
- Files created automatically by the server are owned by the account that runs the server. Files created using `mysql_ssl_rsa_setup` are owned by the user who invoked that program. This can be changed on systems that support the `chown()` system call if the program is invoked by `root` and the `--uid` option is given to specify the user who should own the files.
- On Unix and Unix-like systems, the file access mode is 644 for certificate files (that is, world readable) and 600 for key files (that is, accessible only by the account that runs the server).

To see the contents of an SSL certificate (for example, to check the range of dates over which it is valid), invoke `openssl` directly:

```
shell> openssl x509 -text -in ca.pem
shell> openssl x509 -text -in server-cert.pem
shell> openssl x509 -text -in client-cert.pem
```

It is also possible to check SSL certificate expiration information using this SQL statement:

```
mysql> SHOW STATUS LIKE 'Ssl_server_not%';
+-----------------------+-------------------------+
| Variable_name         | Value                   |
+-----------------------+-------------------------+
| Ssl_server_not_after  | Apr 28 14:16:39 2025 GMT |
| Ssl_server_not_before | May  1 14:16:39 2015 GMT |
+-----------------------+-------------------------+
```

6.4.6.2 Creating SSL Certificates and Keys Using openssl

This section describes how to use the `openssl` command to set up SSL certificate and key files for use by MySQL servers and clients. The first example shows a simplified procedure such as you might use from the command line. The second shows a script that contains more detail. The first two examples are intended for use on Unix and both use the `openssl` command that is part of OpenSSL. The third example describes how to set up SSL files on Windows.

Note

There are easier alternatives to generating the files required for SSL than the procedure described here: Let the server autogenerate them or use the `mysql_ssl_rsa_setup` program. See Section 6.4.6.1, “Creating SSL and RSA Certificates and Keys using MySQL”.

Important

Whatever method you use to generate the certificate and key files, the Common Name value used for the server and client certificates/keys must each differ from the Common Name value used for the CA certificate. Otherwise, the certificate and key files will not work for servers compiled using OpenSSL. A typical error in this case is:

```
ERROR 2026 (HY000): SSL connection error:
error:00000001:lib(0):func(0):reason(1)
```

- Example 1: Creating SSL Files from the Command Line on Unix
- Example 2: Creating SSL Files Using a Script on Unix
- Example 3: Creating SSL Files on Windows

Example 1: Creating SSL Files from the Command Line on Unix

The following example shows a set of commands to create MySQL server and client certificate and key files. You will need to respond to several prompts by the `openssl` commands. To generate test files, you can press Enter to all prompts. To generate files for production use, you should provide nonempty responses.

```
# Create clean environment
shell> rm -rf newcerts
shell> mkdir newcerts && cd newcerts

# Create CA certificate
shell> openssl genrsa 2048 > ca-key.pem
shell> openssl req -new -x509 -nodes -days 3600 \
         -key ca-key.pem -out ca.pem

# Create server certificate, remove passphrase, and sign it
# server-cert.pem = public key, server-key.pem = private key
shell> openssl req -newkey rsa:2048 -days 3600 \
         -nodes -keyout server-key.pem -out server-req.pem
shell> openssl rsa -in server-key.pem -out server-key.pem
shell> openssl x509 -req -in server-req.pem -days 3600 \
         -CA ca.pem -CAkey ca-key.pem -set_serial 01 -out server-cert.pem

# Create client certificate, remove passphrase, and sign it
# client-cert.pem = public key, client-key.pem = private key
shell> openssl req -newkey rsa:2048 -days 3600 \
         -nodes -keyout client-key.pem -out client-req.pem
shell> openssl rsa -in client-key.pem -out client-key.pem
shell> openssl x509 -req -in client-req.pem -days 3600 \
         -CA ca.pem -CAkey ca-key.pem -set_serial 01 -out client-cert.pem
```

After generating the certificates, verify them:

```
shell> openssl verify -CAfile ca.pem server-cert.pem client-cert.pem
server-cert.pem: OK
client-cert.pem: OK
```

To see the contents of a certificate (for example, to check the range of dates over which a certificate is valid), invoke `openssl` like this:

```
shell> openssl x509 -text -in ca.pem
shell> openssl x509 -text -in server-cert.pem
shell> openssl x509 -text -in client-cert.pem
```

Now you have a set of files that can be used as follows:

- `ca.pem`: Use this as the argument to `--ssl-ca` on the server and client sides. (The CA certificate, if used, must be the same on both sides.)
- `server-cert.pem`, `server-key.pem`: Use these as the arguments to `--ssl-cert` and `--ssl-key` on the server side.
- `client-cert.pem`, `client-key.pem`: Use these as the arguments to `--ssl-cert` and `--ssl-key` on the client side.

To use the files for SSL connections, see Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

Example 2: Creating SSL Files Using a Script on Unix

Here is an example script that shows how to set up SSL certificate and key files for MySQL. After executing the script, use the files for SSL connections as described in Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

```
DIR=`pwd`/openssl
PRIV=$DIR/private

mkdir $DIR $PRIV $DIR/newcerts
cp /usr/share/ssl/openssl.cnf $DIR
replace ./demoCA $DIR -- $DIR/openssl.cnf

# Create necessary files: $database, $serial and $new_certs_dir
# directory (optional)

touch $DIR/index.txt
echo "01" > $DIR/serial

#
# Generation of Certificate Authority(CA)
#

openssl req -new -x509 -keyout $PRIV/cakey.pem -out $DIR/ca.pem \
    -days 3600 -config $DIR/openssl.cnf

# Sample output:
# Using configuration from /home/monty/openssl/openssl.cnf
# Generating a 1024 bit RSA private key
# ..............++++++
# .........++++++
# writing new private key to '/home/monty/openssl/private/cakey.pem'
# Enter PEM pass phrase:
# Verifying password - Enter PEM pass phrase:
# -----
# You are about to be asked to enter information that will be
# incorporated into your certificate request.
# What you are about to enter is what is called a Distinguished Name
# or a DN.
# There are quite a few fields but you can leave some blank
# For some fields there will be a default value,
# If you enter '.', the field will be left blank.
# -----
# Country Name (2 letter code) [AU]:FI
# State or Province Name (full name) [Some-State]:.
# Locality Name (eg, city) []:
# Organization Name (eg, company) [Internet Widgits Pty Ltd]:MySQL AB
# Organizational Unit Name (eg, section) []:
# Common Name (eg, YOUR name) []:MySQL admin
# Email Address []:

#
# Create server request and key
#
openssl req -new -keyout $DIR/server-key.pem -out \
    $DIR/server-req.pem -days 3600 -config $DIR/openssl.cnf
```

```
# Sample output:
# Using configuration from /home/monty/openssl/openssl.cnf
# Generating a 1024 bit RSA private key
# ..++++++
# ..........++++++
# writing new private key to '/home/monty/openssl/server-key.pem'
# Enter PEM pass phrase:
# Verifying password - Enter PEM pass phrase:
# -----
# You are about to be asked to enter information that will be
# incorporated into your certificate request.
# What you are about to enter is what is called a Distinguished Name
# or a DN.
# There are quite a few fields but you can leave some blank
# For some fields there will be a default value,
# If you enter '.', the field will be left blank.
# -----
# Country Name (2 letter code) [AU]:FI
# State or Province Name (full name) [Some-State]:.
# Locality Name (eg, city) []:
# Organization Name (eg, company) [Internet Widgits Pty Ltd]:MySQL AB
# Organizational Unit Name (eg, section) []:
# Common Name (eg, YOUR name) []:MySQL server
# Email Address []:
#
# Please enter the following 'extra' attributes
# to be sent with your certificate request
# A challenge password []:
# An optional company name []:

#
# Remove the passphrase from the key
#
openssl rsa -in $DIR/server-key.pem -out $DIR/server-key.pem

#
# Sign server cert
#
openssl ca -cert $DIR/ca.pem -policy policy_anything \
    -out $DIR/server-cert.pem -config $DIR/openssl.cnf \
    -infiles $DIR/server-req.pem

# Sample output:
# Using configuration from /home/monty/openssl/openssl.cnf
# Enter PEM pass phrase:
# Check that the request matches the signature
# Signature ok
# The Subjects Distinguished Name is as follows
# countryName           :PRINTABLE:'FI'
# organizationName      :PRINTABLE:'MySQL AB'
# commonName            :PRINTABLE:'MySQL admin'
# Certificate is to be certified until Sep 13 14:22:46 2003 GMT
# (365 days)
# Sign the certificate? [y/n]:y
#
#
# 1 out of 1 certificate requests certified, commit? [y/n]y
# Write out database with 1 new entries
# Data Base Updated

#
# Create client request and key
#
openssl req -new -keyout $DIR/client-key.pem -out \
    $DIR/client-req.pem -days 3600 -config $DIR/openssl.cnf

# Sample output:
# Using configuration from /home/monty/openssl/openssl.cnf
# Generating a 1024 bit RSA private key
# ....................................++++++
# ..........................................++++++
# writing new private key to '/home/monty/openssl/client-key.pem'
```

```
# Enter PEM pass phrase:
# Verifying password - Enter PEM pass phrase:
# -----
# You are about to be asked to enter information that will be
# incorporated into your certificate request.
# What you are about to enter is what is called a Distinguished Name
# or a DN.
# There are quite a few fields but you can leave some blank
# For some fields there will be a default value,
# If you enter '.', the field will be left blank.
# -----
# Country Name (2 letter code) [AU]:FI
# State or Province Name (full name) [Some-State]:.
# Locality Name (eg, city) []:
# Organization Name (eg, company) [Internet Widgits Pty Ltd]:MySQL AB
# Organizational Unit Name (eg, section) []:
# Common Name (eg, YOUR name) []:MySQL user
# Email Address []:
#
# Please enter the following 'extra' attributes
# to be sent with your certificate request
# A challenge password []:
# An optional company name []:

#
# Remove the passphrase from the key
#
openssl rsa -in $DIR/client-key.pem -out $DIR/client-key.pem

#
# Sign client cert
#

openssl ca -cert $DIR/ca.pem -policy policy_anything \
    -out $DIR/client-cert.pem -config $DIR/openssl.cnf \
    -infiles $DIR/client-req.pem

# Sample output:
# Using configuration from /home/monty/openssl/openssl.cnf
# Enter PEM pass phrase:
# Check that the request matches the signature
# Signature ok
# The Subjects Distinguished Name is as follows
# countryName           :PRINTABLE:'FI'
# organizationName      :PRINTABLE:'MySQL AB'
# commonName            :PRINTABLE:'MySQL user'
# Certificate is to be certified until Sep 13 16:45:17 2003 GMT
# (365 days)
# Sign the certificate? [y/n]:y
#
#
# 1 out of 1 certificate requests certified, commit? [y/n]y
# Write out database with 1 new entries
# Data Base Updated

#
# Create a my.cnf file that you can use to test the certificates
#

cat <<EOF > $DIR/my.cnf
[client]
ssl-ca=$DIR/ca.pem
ssl-cert=$DIR/client-cert.pem
ssl-key=$DIR/client-key.pem
[mysqld]
ssl-ca=$DIR/ca.pem
ssl-cert=$DIR/server-cert.pem
ssl-key=$DIR/server-key.pem
EOF
```

Example 3: Creating SSL Files on Windows

Download OpenSSL for Windows if it is not installed on your system. An overview of available packages can be seen here:

```
http://www.slproweb.com/products/Win32OpenSSL.html
```

Choose the Win32 OpenSSL Light or Win64 OpenSSL Light package, depending on your architecture (32-bit or 64-bit). The default installation location will be `C:\OpenSSL-Win32` or `C:\OpenSSL-Win64`, depending on which package you downloaded. The following instructions assume a default location of `C:\OpenSSL-Win32`. Modify this as necessary if you are using the 64-bit package.

If a message occurs during setup indicating `'...critical component is missing: Microsoft Visual C++ 2008 Redistributables'`, cancel the setup and download one of the following packages as well, again depending on your architecture (32-bit or 64-bit):

- Visual C++ 2008 Redistributables (x86), available at:

  ```
  http://www.microsoft.com/downloads/details.aspx?familyid=9B2DA534-3E03-4391-8A4D-074B9F2BC1BF
  ```

- Visual C++ 2008 Redistributables (x64), available at:

  ```
  http://www.microsoft.com/downloads/details.aspx?familyid=bd2a6171-e2d6-4230-b809-9a8d7548c1b6
  ```

After installing the additional package, restart the OpenSSL setup procedure.

During installation, leave the default `C:\OpenSSL-Win32` as the install path, and also leave the default option `'Copy OpenSSL DLL files to the Windows system directory'` selected.

When the installation has finished, add `C:\OpenSSL-Win32\bin` to the Windows System Path variable of your server:

1. On the Windows desktop, right-click the **My Computer** icon, and select **Properties**.
2. Select the **Advanced** tab from the **System Properties** menu that appears, and click the **Environment Variables** button.
3. Under **System Variables**, select **Path**, then click the **Edit** button. The **Edit System Variable** dialogue should appear.
4. Add `';C:\OpenSSL-Win32\bin'` to the end (notice the semicolon).
5. Press OK 3 times.
6. Check that OpenSSL was correctly integrated into the Path variable by opening a new command console (`Start>Run>cmd.exe`) and verifying that OpenSSL is available:

   ```
   Microsoft Windows [Version ...]
   Copyright (c) 2006 Microsoft Corporation. All rights reserved.

   C:\Windows\system32>cd \

   C:\>openssl
   OpenSSL> exit <<< If you see the OpenSSL prompt, installation was successful.

   C:\>
   ```

Depending on your version of Windows, the preceding path-setting instructions might differ slightly.

After OpenSSL has been installed, use instructions similar to those from Example 1 (shown earlier in this section), with the following changes:

- Change the following Unix commands:

```
# Create clean environment
shell> rm -rf newcerts
shell> mkdir newcerts && cd newcerts
```

  On Windows, use these commands instead:

```
# Create clean environment
C:\> md c:\newcerts
C:\> cd c:\newcerts
```

- When a `'\'` character is shown at the end of a command line, this `'\'` character must be removed and the command lines entered all on a single line.

After generating the certificate and key files, to use them for SSL connections, see Section 6.4.4, "Configuring MySQL to Use Secure Connections".

6.4.6.3 Creating RSA Keys Using openssl

This section describes how to use the `openssl` command to set up the RSA key files that enable MySQL to support secure password exchange over unencrypted connections for accounts authenticated by the `sha256_password` plugin.

Note

There are easier alternatives to generating the files required for RSA than the procedure described here: Let the server autogenerate them or use the `mysql_ssl_rsa_setup` program. See Section 6.4.6.1, "Creating SSL and RSA Certificates and Keys using MySQL".

To create the RSA private and public key-pair files, run these commands while logged into the system account used to run the MySQL server so the files will be owned by that account:

```
openssl genrsa -out private_key.pem 2048
openssl rsa -in private_key.pem -pubout -out public_key.pem
```

Those commands create 2,048-bit keys. To create stronger keys, use a larger value.

Then set the access modes for the key files. The private key should be readable only by the server, whereas the public key can be freely distributed to client users:

```
chmod 400 private_key.pem
chmod 444 public_key.pem
```

6.4.7 Connecting to MySQL Remotely from Windows with SSH

This section describes how to get a secure connection to a remote MySQL server with SSH. The information was provided by David Carlson `<dcarlson@mplcomm.com>`.

1. Install an SSH client on your Windows machine. For a comparison of SSH clients, see http://en.wikipedia.org/wiki/Comparison_of_SSH_clients.

2. Start your Windows SSH client. Set `Host_Name = yourmysqlserver_URL_or_IP`. Set `userid=your_userid` to log in to your server. This `userid` value might not be the same as the user name of your MySQL account.

3. Set up port forwarding. Either do a remote forward (Set `local_port: 3306`, `remote_host: yourmysqlservername_or_ip`, `remote_port: 3306`) or a local forward (Set `port: 3306`, `host: localhost`, `remote port: 3306`).

4. Save everything, otherwise you will have to redo it the next time.
5. Log in to your server with the SSH session you just created.
6. On your Windows machine, start some ODBC application (such as Access).
7. Create a new file in Windows and link to MySQL using the ODBC driver the same way you normally do, except type in `localhost` for the MySQL host server, not *`yourmysqlservername`*.

At this point, you should have an ODBC connection to MySQL, encrypted using SSH.

6.5 Security Plugins

MySQL includes several plugins that implement security features:

- Plugins for authenticating attempts by clients to connect to MySQL Server. Plugins are available for several authentication protocols. For general discussion of the authentication process, see Section 6.3.10, “Pluggable Authentication”. For characteristics of specific authentication plugins, see Section 6.5.1, “Authentication Plugins”.
- A password-validation plugin for implementing password strength policies and assessing the strength of potential passwords. See Section 6.5.3, “The Password Validation Plugin”.
- A keyring plugin that provides secure storage for sensitive information. See Section 6.5.4, “The MySQL Keyring”.

6.5.1 Authentication Plugins

The following sections describe pluggable authentication methods available in MySQL and the plugins that implement these methods. For general discussion of the authentication process, see Section 6.3.10, “Pluggable Authentication”.

The default plugin is `mysql_native_password` unless the `default_authentication_plugin` system variable is set otherwise.

6.5.1.1 Native Pluggable Authentication

MySQL includes a `mysql_native_password` plugin that implements native authentication; that is, authentication based on the password hashing method in use from before the introduction of pluggable authentication.

The following table shows the plugin names on the server and client sides.

Table 6.10 Plugin and Library Names for Native Password Authentication

Server-side plugin name	`mysql_native_password`
Client-side plugin name	`mysql_native_password`
Library file name	None (plugins are built in)

The following sections provide installation and usage information specific to native pluggable authentication:

- Installing Native Pluggable Authentication
- Using Native Pluggable Authentication

For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”.

Installing Native Pluggable Authentication

The `mysql_native_password` plugin exists in server and client forms:

- The server-side plugin is built into the server, need not be loaded explicitly, and cannot be disabled by unloading it.
- The client-side plugin is built into the `libmysqlclient` client library and available to any program linked against `libmysqlclient`.

Using Native Pluggable Authentication

MySQL client programs use `mysql_native_password` by default. The `--default-auth` option can be used as a hint about which client-side plugin the program can expect to use:

```
shell> mysql --default-auth=mysql_native_password ...
```

6.5.1.2 SHA-256 Pluggable Authentication

MySQL provides an authentication plugin that implements SHA-256 hashing for user account passwords.

Important

To connect to the server using an account that authenticates with the `sha256_password` plugin, you must use either an SSL connection or an unencrypted connection that encrypts the password using RSA, as described later in this section. Either way, the `sha256_password` plugin uses MySQL's SSL capabilities. See Section 6.4, “Using Secure Connections”.

The following table shows the plugin names on the server and client sides.

Table 6.11 Plugin and Library Names for SHA-256 Authentication

Server-side plugin name	`sha256_password`
Client-side plugin name	`sha256_password`
Library file name	None (plugins are built in)

The following sections provide installation and usage information specific to SHA-256 pluggable authentication:

- Installing SHA-256 Pluggable Authentication
- Using sha256 Pluggable Authentication

For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”.

Installing SHA-256 Pluggable Authentication

The `sha256_password` plugin exists in server and client forms:

- The server-side plugin is built into the server, need not be loaded explicitly, and cannot be disabled by unloading it.
- The client-side plugin is built into the `libmysqlclient` client library and available to any program linked against `libmysqlclient`.

Using sha256 Pluggable Authentication

To set up an account that uses the `sha256_password` plugin for SHA-256 password hashing, use the following statement:

```
CREATE USER 'sha256user'@'localhost'
IDENTIFIED WITH sha256_password BY 'Sh@256Pa33';
```

Alternatively, start the server with the default authentication plugin set to `sha256_password`. For example, put these lines in the server option file:

```
[mysqld]
default_authentication_plugin=sha256_password
```

That causes the `sha256_password` plugin to be used by default for new accounts. As a result, it is possible to create the account and set its password without naming the plugin explicitly using this `CREATE USER` syntax:

```
CREATE USER 'sha256user'@'localhost' IDENTIFIED BY 'Sh@256Pa33';
```

In this case, the server assigns the `sha256_password` plugin to the account and encrypts the password using SHA-256.

Accounts in the `mysql.user` table that use SHA-256 passwords can be identified as rows with `'sha256_password'` in the `plugin` column and a SHA-256 password hash in the `authentication_string` column.

Another consequence of using `sha256_password` as the default authentication plugin is that to create an account that uses a different plugin, you must specify that plugin using an `IDENTIFIED WITH` clause in the `CREATE USER` statement. For example, to use the `mysql_native_password` plugin, use this statement:

```
CREATE USER 'nativeuser'@'localhost'
IDENTIFIED WITH mysql_native_password BY 'N@tivePa33';
```

MySQL can be compiled using either OpenSSL or yaSSL (see Section 6.4.1, “OpenSSL Versus yaSSL”). The `sha256_password` plugin works with distributions compiled using either package, but if MySQL is compiled using OpenSSL, RSA encryption is available and `sha256_password` implements the following additional capabilities. (To enable these capabilities, you must also follow the RSA configuration procedure given later in this section.)

- It is possible for the client to transmit passwords to the server using RSA encryption during the client connection process, as described later.
- The server exposes two additional system variables, `sha256_password_private_key_path` and `sha256_password_public_key_path`. It is intended that the database administrator will set these to the names of the RSA private and public key-pair files at server startup if the key files have names that differ from the system variable default values.
- The server exposes a status variable, `Rsa_public_key`, that displays the RSA public key value.
- The `mysql` and `mysqltest` client programs support a `--server-public-key-path` option for specifying an RSA public key file explicitly.

For clients that use the `sha256_password` plugin, passwords are never exposed as cleartext when connecting to the server. How password transmission occurs depends on whether an SSL connection is used and whether RSA encryption is available:

- If an SSL connection is used, the password is sent as cleartext but cannot be snooped because the connection is encrypted using SSL.

- If an SSL connection is not used but RSA encryption is available, the password is sent within an unencrypted connection, but the password is RSA-encrypted to prevent snooping. When the server receives the password, it decrypts it. A scramble is used in the encryption to prevent repeat attacks.
- If an SSL connection is not used and RSA encryption is not available, the `sha256_password` plugin causes the connection attempt to fail because the password cannot be sent without being exposed as cleartext.

As mentioned previously, RSA password encryption is available only if MySQL was compiled using OpenSSL. The implication for MySQL distributions compiled using yaSSL is that SHA-256 passwords can be used only when clients use SSL to access the server. See Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

Assuming that MySQL has been compiled using OpenSSL, the following procedure describes how to enable RSA encryption of passwords during the client connection process:

1. Create the RSA private and public key-pair files using the instructions in Section 6.4.6, “Creating SSL and RSA Certificates and Keys”.
2. If the private and public key files are located in the data directory and are named `private_key.pem` and `public_key.pem` (the default values of the `sha256_password_private_key_path` and `sha256_password_public_key_path` system variables), the server uses them automatically at startup.

   Otherwise, in the server option file, set the system variables to the key file names. If the files are located in the server data directory, you need not specify their full path names:

   ```
   [mysqld]
   sha256_password_private_key_path=myprivkey.pem
   sha256_password_public_key_path=mypubkey.pem
   ```

   If the key files are not located in the data directory, or to make their locations explicit in the system variable values, use full path names:

   ```
   [mysqld]
   sha256_password_private_key_path=/usr/local/mysql/myprivkey.pem
   sha256_password_public_key_path=/usr/local/mysql/mypubkey.pem
   ```

3. Restart the server, then connect to it and check the `Rsa_public_key` status variable value. The value will differ from that shown here, but should be nonempty:

   ```
   mysql> SHOW STATUS LIKE 'Rsa_public_key'\G
   *************************** 1. row ***************************
   Variable_name: Rsa_public_key
           Value: -----BEGIN PUBLIC KEY-----
   MIGfMA0GCSqGSIb3DQEBAQUAA4GNADCBiQKBgQDO9nRUDd+KvSZgY7cNBZMNpwX6
   MvE1PbJFXO7u18nJ9lwc99Du/E7lw6CVXw7VKrXPeHbVQUzGyUNkf45Nz/ckaaJa
   aLgJOBCIDmNVnyU54OT/1lcs2xiyfaDMe8fCJ64ZwTnKbY2gkt1IMjUAB5Ogd5kJ
   g8aV7EtKwyhHb0c30QIDAQAB
   -----END PUBLIC KEY-----
   ```

   If the value is empty, the server found some problem with the key files. Check the error log for diagnostic information.

After the server has been configured with the RSA key files, accounts that authenticate with the `sha256_password` plugin have the option of using those key files to connect to the server. As mentioned previously, such accounts can use either an SSL connection (in which case RSA is not used) or an unencrypted connection that encrypts the password using RSA. Assume for the following discussion that SSL is not used. Connecting to the server involves no special preparation on the client side. For example:

```
shell> mysql --ssl-mode=DISABLED -u sha256user -p
Enter password: Sh@256Pa33
```

For connection attempts by `sha256user`, the server determines that `sha256_password` is the appropriate authentication plugin and invokes it. The plugin finds that the connection does not use SSL and thus requires the password to be transmitted using RSA encryption. In this case, the plugin sends the RSA public key to the client, which uses it to encrypt the password and returns the result to the server. The plugin uses the RSA key on the server side to decrypt the password and accepts or rejects the connection based on whether the password is correct.

The server sends the public key to the client as needed, but if a copy of the RSA public key is available on the client host, the client can use it to save a round trip in the client/server protocol:

```
shell> mysql --ssl-mode=DISABLED -u sha256user -p --server-public-key-path=file_name
Enter password: Sh@256Pa33
```

The public key value in the file named by the `--server-public-key-path` option should be the same as the key value in the server-side file named by the `sha256_password_public_key_path` system variable. If the key file contains a valid public key value but the value is incorrect, an access-denied error occurs. If the key file does not contain a valid public key, the client program cannot use it. In this case, the `sha256_password` plugin sends the public key to the client as if no `--server-public-key-path` option had been specified.

Client users can get the RSA public key two ways:

- The database administrator can provide a copy of the public key file.
- A client user who can connect to the server some other way can use a `SHOW STATUS LIKE 'Rsa_public_key'` statement and save the returned key value in a file.

6.5.1.3 Client-Side Cleartext Pluggable Authentication

A client-side authentication plugin is available that sends the password to the server without hashing or encryption. This plugin is built into the MySQL client library.

The following table shows the plugin name.

Table 6.12 Plugin and Library Names for Cleartext Authentication

Server-side plugin name	None, see discussion
Client-side plugin name	`mysql_clear_password`
Library file name	None (plugin is built in)

With many MySQL authentication methods, the client performs hashing or encryption of the password before sending it to the server. This enables the client to avoid sending the password in clear text.

Hashing or encryption cannot be done for authentication schemes that require the server to receive the password as entered on the client side. In such cases, the client-side `mysql_clear_password` plugin is used to send the password to the server in clear text. There is no corresponding server-side plugin. Rather, the client-side plugin can be used by any server-side plugin that needs a cleartext password. (The PAM authentication plugin is one such; see PAM Pluggable Authentication.)

The following discussion provides usage information specific to clear text pluggable authentication. For For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”.

Note

Sending passwords in clear text may be a security problem in some configurations. To avoid problems if there is any possibility that the password would be intercepted, clients should connect to MySQL Server using a method

that protects the password. Possibilities include SSL (see Section 6.4, “Using Secure Connections”), IPsec, or a private network.

To make inadvertent use of the `mysql_clear_password` plugin less likely, clients must explicitly enable it. This can be done several ways:

- Set the `LIBMYSQL_ENABLE_CLEARTEXT_PLUGIN` environment variable to a value that begins with `1`, `Y`, or `y`. This enables the plugin for all client connections.
- The `mysql`, `mysqladmin`, `mysqlcheck`, `mysqldump`, `mysqlshow`, and `mysqlslap` client programs support an `--enable-cleartext-plugin` option that enables the plugin on a per-invocation basis.
- The `mysql_options()` C API function supports a `MYSQL_ENABLE_CLEARTEXT_PLUGIN` option that enables the plugin on a per-connection basis. Also, any program that uses `libmysqlclient` and reads option files can enable the plugin by including an `enable-cleartext-plugin` option in an option group read by the client library.

6.5.1.4 No-Login Pluggable Authentication

The `mysql_no_login` server-side authentication plugin prevents all client connections to any account that uses it. Use cases for such a plugin includes accounts that must be able to execute stored programs and views with elevated privileges without exposing those privileges to ordinary users, and proxied accounts that should never permit direct login but are accessed only through proxy accounts.

The following table shows the plugin and library file names. The file name suffix might differ on your system. The file must be located in the directory named by the `plugin_dir` system variable.

Table 6.13 Plugin and Library Names for “No Login” Authentication

Server-side plugin name	`mysql_no_login`
Client-side plugin name	None
Library file name	`mysql_no_login.so`

The following sections provide installation and usage information specific to no-login pluggable authentication:

- Installing No-Login Pluggable Authentication
- Uninstalling No-Login Pluggable Authentication
- Using No-Login Pluggable Authentication

For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”. For proxy user information, see Section 6.3.11, “Proxy Users”.

Installing No-Login Pluggable Authentication

This section describes how to install the no-login authentication plugin. For general information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base name is `mysql_no_login`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To load the plugin at server startup, use the `--plugin-load-add` option to name the library file that contains it. With this plugin-loading method, the option must be given each time the server starts. For example, put these lines in your `my.cnf` file (adjust the `.so` suffix for your platform as necessary):

```
[mysqld]
plugin-load-add=mysql_no_login.so
```

Alternatively, to register the plugin at runtime, use this statement (adjust the extension as necessary):

```
INSTALL PLUGIN mysql_no_login SONAME 'mysql_no_login.so';
```

`INSTALL PLUGIN` loads a plugin, and also registers it in the `mysql.plugins` table to cause the plugin to be loaded for each subsequent normal server startup.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, "Obtaining Server Plugin Information"). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE '%login%';
+---------------+---------------+
| PLUGIN_NAME    | PLUGIN_STATUS |
+---------------+---------------+
| mysql_no_login | ACTIVE        |
+---------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

To associate a MySQL account with the no-login plugin, use the plugin name `mysql_no_login` in the `IDENTIFIED WITH` clause of the `CREATE USER` statement that creates the account. (See Using No-Login Pluggable Authentication.)

Uninstalling No-Login Pluggable Authentication

The method used to uninstall the no-login authentication plugin depends on how you installed it:

- If you installed the plugin at server startup using a `--plugin-load-add` option, restart the server without the option.
- If you installed the plugin at runtime using `INSTALL PLUGIN`, it remains installed across server restarts. To uninstall it, use `UNINSTALL PLUGIN`:

  ```
  UNINSTALL PLUGIN mysql_no_login;
  ```

Using No-Login Pluggable Authentication

This section describes how to use the no-login authentication plugin to prevent connections from MySQL client programs to the server. It is assumed that the server is running with the server-side plugin enabled, as described in Installing No-Login Pluggable Authentication.

To refer to the no-login authentication plugin in the `IDENTIFIED WITH` clause of a `CREATE USER` statement, use the name `mysql_no_login`.

An account that authenticates using `mysql_no_login` may be used as the `DEFINER` for stored program and view objects. If such an object definition also includes `SQL SECURITY DEFINER`, it executes with that account's privileges. DBAs can use this behavior to provide access to confidential or sensitive data that is exposed only through well-controlled interfaces.

The following example provides a simple illustration of these principles. It defines an account that does not permit client connections, and associates with it a view that exposes only certain columns of the `mysql.user` table:

```
CREATE DATABASE nologindb;
CREATE USER 'nologin'@'localhost'
  IDENTIFIED WITH mysql_no_login;
GRANT ALL ON nologindb.*
  TO 'nologin'@'localhost';
GRANT SELECT ON mysql.user
  TO 'nologin'@'localhost';
CREATE DEFINER = 'nologin'@'localhost'
  SQL SECURITY DEFINER
  VIEW nologindb.myview
  AS SELECT User, Host FROM mysql.user;
```

To provide protected access to the view to an ordinary user, do this:

```
GRANT SELECT ON nologindb.myview
  TO 'ordinaryuser'@'localhost';
```

Now the ordinary user can use the view to access the limited information it presents:

```
SELECT * FROM nologindb.myview;
```

Attempts by the user to access columns other than those exposed by the view result in an error, as do all attempts to select from the view by users not granted access to it.

Note

Because the `nologin` account cannot be used directly, the operations required to set up objects that it uses must be performed by `root` or similar account with the privileges required to create the objects and set `DEFINER` values.

An account that authenticates using `mysql_no_login` may be used as a proxied base user for proxy accounts:

```
-- create proxied account
CREATE USER 'proxy_base'@'localhost'
  IDENTIFIED WITH mysql_no_login;
-- grant privileges to proxied account
GRANT ...  TO 'proxy_base'@'localhost';
- permit real_user to be proxy for proxied account
GRANT PROXY ON 'proxy_base'@'localhost'
  TO 'real_user'@'localhost';
```

This enables clients to access MySQL through the proxy account (`real_user`) but not to bypass the proxy mechanism by connecting directly as the proxied user (`proxy_base`).

6.5.1.5 Socket Peer-Credential Pluggable Authentication

The server-side `auth_socket` authentication plugin authenticates clients that connect from the local host through the Unix socket file. The plugin uses the `SO_PEERCRED` socket option to obtain information about the user running the client program. Thus, the plugin can be used only on systems that support the `SO_PEERCRED` option, such as Linux.

The source code for this plugin can be examined as a relatively simple example demonstrating how to write a loadable authentication plugin.

The following table shows the plugin and library file names. The file must be located in the directory named by the `plugin_dir` system variable.

Table 6.14 Plugin and Library Names for Socket Peer-Credential Authentication

Server-side plugin name	`auth_socket`

Client-side plugin name	None, see discussion
Library file name	`auth_socket.so`

The following sections provide installation and usage information specific to socket pluggable authentication:

- Installing Socket Pluggable Authentication
- Uninstalling Socket Pluggable Authentication
- Using Socket Pluggable Authentication

For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”.

Installing Socket Pluggable Authentication

This section describes how to install the socket authentication plugin. For general information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

To load the plugin at server startup, use the `--plugin-load-add` option to name the library file that contains it. With this plugin-loading method, the option must be given each time the server starts. For example, put these lines in your `my.cnf` file:

```
[mysqld]
plugin-load-add=auth_socket.so
```

Alternatively, to register the plugin at runtime, use this statement (adjust the extension as necessary):

```
INSTALL PLUGIN auth_socket SONAME 'auth_socket.so';
```

`INSTALL PLUGIN` loads a plugin, and also registers it in the `mysql.plugins` table to cause the plugin to be loaded for each subsequent normal server startup.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE '%socket%';
+-------------+---------------+
| PLUGIN_NAME | PLUGIN_STATUS |
+-------------+---------------+
| auth_socket | ACTIVE        |
+-------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

To associate a MySQL account with the socket plugin, use the plugin name `auth_socket` in the `IDENTIFIED WITH` clause of the `CREATE USER` statement that creates the account. (See Using Socket Pluggable Authentication.)

Uninstalling Socket Pluggable Authentication

The method used to uninstall the socket authentication plugin depends on how you installed it:

- If you installed the plugin at server startup using a `--plugin-load-add` option, restart the server without the option.
- If you installed the plugin at runtime using `INSTALL PLUGIN`, it remains installed across server restarts. To uninstall it, use `UNINSTALL PLUGIN`:

```
UNINSTALL PLUGIN auth_socket;
```

Using Socket Pluggable Authentication

The socket plugin checks whether the socket user name matches the MySQL user name specified by the client program to the server. If the names do not match, the plugin also checks whether the socket user name matches the name specified in the `authentication_string` column of the `mysql.user` table row. If a match is found, the plugin permits the connection.

Suppose that a MySQL account is created for a user named `valerie` who is to be authenticated by the `auth_socket` plugin for connections from the local host through the socket file:

```
CREATE USER 'valerie'@'localhost' IDENTIFIED WITH auth_socket;
```

If a user on the local host with a login name of `stefanie` invokes `mysql` with the option `--user=valerie` to connect through the socket file, the server uses `auth_socket` to authenticate the client. The plugin determines that the `--user` option value (`valerie`) differs from the client user's name (`stephanie`) and refuses the connection. If a user named `valerie` tries the same thing, the plugin finds that the user name and the MySQL user name are both `valerie` and permits the connection. However, the plugin refuses the connection even for `valerie` if the connection is made using a different protocol, such as TCP/IP.

6.5.1.6 Test Pluggable Authentication

MySQL includes a test plugin that checks account credentials and logs success or failure to the server error log. This is a loadable plugin (not built in) and must be installed prior to use.

The test plugin source code is separate from the server source, unlike the built-in native plugin, so it can be examined as a relatively simple example demonstrating how to write a loadable authentication plugin.

Note

This plugin is intended for testing and development purposes, and is not for use in production environments or on servers that are exposed to public networks.

The following table shows the plugin and library file names. The file name suffix might differ on your system. The file must be located in the directory named by the `plugin_dir` system variable.

Table 6.15 Plugin and Library Names for Test Authentication

Server-side plugin name	`test_plugin_server`
Client-side plugin name	`auth_test_plugin`
Library file name	`auth_test_plugin.so`

The following sections provide installation and usage information specific to test pluggable authentication:

- Installing Test Pluggable Authentication
- Uninstalling Test Pluggable Authentication
- Using Test Pluggable Authentication

For general information about pluggable authentication in MySQL, see Section 6.3.10, “Pluggable Authentication”.

Installing Test Pluggable Authentication

This section describes how to install the test authentication plugin. For general information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

To load the plugin at server startup, use the `--plugin-load-add` option to name the library file that contains it. With this plugin-loading method, the option must be given each time the server starts. For example, put these lines in your `my.cnf` file:

```
[mysqld]
plugin-load-add=auth_test_plugin.so
```

Alternatively, to register the plugin at runtime, use this statement (adjust the extension as necessary):

```
INSTALL PLUGIN test_plugin_server SONAME 'auth_test_plugin.so';
```

`INSTALL PLUGIN` loads a plugin, and also registers it in the `mysql.plugins` table to cause the plugin to be loaded for each subsequent normal server startup.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE '%test_plugin%';
+--------------------+---------------+
| PLUGIN_NAME        | PLUGIN_STATUS |
+--------------------+---------------+
| test_plugin_server | ACTIVE        |
+--------------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

To associate a MySQL account with the test plugin, use the plugin name `test_plugin_server` in the `IDENTIFIED WITH` clause of the `CREATE USER` statement that creates the account. (See Using Test Pluggable Authentication.)

Uninstalling Test Pluggable Authentication

The method used to uninstall the test authentication plugin depends on how you installed it:

- If you installed the plugin at server startup using a `--plugin-load-add` option, restart the server without the option.
- If you installed the plugin at runtime using `INSTALL PLUGIN`, it remains installed across server restarts. To uninstall it, use `UNINSTALL PLUGIN`:

  ```
  UNINSTALL PLUGIN test_plugin_server;
  ```

Using Test Pluggable Authentication

To use the test authentication plugin, create an account and name that plugin in the `IDENTIFIED WITH` clause:

```
CREATE USER 'testuser'@'localhost'
IDENTIFIED WITH test_plugin_server
BY 'testpassword';
```

Then provide the `--user` and `--password` options for that account when you connect to the server. For example:

```
shell> mysql --user=testuser --password
Enter password: testpassword
```

The plugin fetches the password as received from the client and compares it with the value stored in the `authentication_string` column of the account row in the `mysql.user` table. If the two values match, the plugin returns the `authentication_string` value as the new effective user ID.

You can look in the server error log for a message indicating whether authentication succeeded (notice that the password is reported as the "user"):

```
[Note] Plugin test_plugin_server reported:
'successfully authenticated user testpassword'
```

6.5.2 The Connection-Control Plugins

MySQL Server includes a plugin library that enables administrators to introduce an increasing delay in server response to clients after a certain number of consecutive failed connection attempts. This capability provides a deterrent that slows down brute force attacks that attempt to access MySQL user accounts. The plugin library contains two plugins:

- `CONNECTION_CONTROL` checks incoming connections and adds a delay to server responses as necessary. This plugin also exposes system variables that enable plugin operation to be configured and a status variable that provides rudimentary monitoring information.

  The `CONNECTION_CONTROL` plugin uses the audit plugin interface (see Section 28.2.4.8, "Writing Audit Plugins"). To collect information, it subscribes to the `MYSQL_AUDIT_CONNECTION_CLASSMASK` event class, and processes `MYSQL_AUDIT_CONNECTION_CONNECT` and `MYSQL_AUDIT_CONNECTION_CHANGE_USER` subevents to check whether the server should introduce a delay before responding to client connection attempts.

- `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` implements an `INFORMATION_SCHEMA` table that exposes more detailed monitoring information for failed connection attempts.

The following sections provide information about connection-control plugin installation and configuration. For information about the `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` table, see Section 24.32.1, "The INFORMATION_SCHEMA CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS Table".

6.5.2.1 Connection-Control Plugin Installation

This section describes how to install the connection-control plugins, `CONNECTION_CONTROL` and `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS`. For general information about installing plugins, see Section 5.6.2, "Installing and Uninstalling Plugins".

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base name is `connection_control`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To load the plugins at server startup, use the `--plugin-load-add` option to name the library file that contains them. With this plugin-loading method, the option must be given each time the server starts. For example, put these lines in your `my.cnf` file (adjust the `.so` suffix for your platform as necessary):

```
[mysqld]
plugin-load-add=connection_control.so
```

Alternatively, to register the plugins at runtime, use these statements (adjust the extension as necessary):

```
INSTALL PLUGIN CONNECTION_CONTROL SONAME 'connection_control.so';
INSTALL PLUGIN CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS SONAME 'connection_control.so';
```

`INSTALL PLUGIN` loads a plugin, and also registers it in the `mysql.plugins` table to cause the plugin to be loaded for each subsequent normal server startup.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE 'connection%';
+---------------------------------------+---------------+
| PLUGIN_NAME                           | PLUGIN_STATUS |
+---------------------------------------+---------------+
| CONNECTION_CONTROL                    | ACTIVE        |
| CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS | ACTIVE        |
+---------------------------------------+---------------+
```

If a plugin fails to initialize, check the server error log for diagnostic messages.

If the plugins have been previously registered with `INSTALL PLUGIN` or are loaded with `--plugin-load-add`, you can use the `--connection-control` and `--connection-control-failed-login-attempts` options at server startup to control plugin activation. For example, to load the plugins at startup and prevent them from being removed at runtime, use these options:

```
[mysqld]
plugin-load-add=connection_control.so
connection-control=FORCE_PLUS_PERMANENT
connection-control-failed-login-attempts=FORCE_PLUS_PERMANENT
```

If it is desired to prevent the server from running without a given connection-control plugin, use an option value of `FORCE` or `FORCE_PLUS_PERMANENT` to force server startup to fail if the plugin does not initialize successfully.

Note

It is possible to install one plugin without the other, but both must be installed for full connection-control capability. In particular, installing only the `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` plugin is of little use because without the `CONNECTION_CONTROL` plugin to provide the data that populates the `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` table, retrievals from the table will always be empty.

Connection Delay Configuration

To enable you to configure its operation, the `CONNECTION_CONTROL` plugin exposes several system variables:

- `connection_control_failed_connections_threshold`: The number of consecutive failed connection attempts permitted to clients before the server adds a delay for subsequent connection attempts.

- `connection_control_min_connection_delay`: The amount of delay to add for each consecutive connection failure above the threshold.
- `connection_control_max_connection_delay`: The maximum delay to add.

To entirely disable checking for failed connection attempts, set `connection_control_failed_connections_threshold` to zero. If `connection_control_failed_connections_threshold` is nonzero, the amount of delay is zero up through that many consecutive failed connection attempts. Thereafter, the amount of delay is the number of failed attempts above the threshold, multiplied by `connection_control_min_connection_delay` milliseconds. For example, with the default `connection_control_failed_connections_threshold` and `connection_control_min_connection_delay` values of 3 and 1000, respectively, there is no delay for the first three consecutive failed connection attempts by a client, a delay of 1000 milliseconds for the fourth failed attempt, 2000 milliseconds for the fifth failed attempt, and so on, up to the maximum delay permitted by `connection_control_max_connection_delay`.

You can set the `CONNECTION_CONTROL` system variables at server startup or runtime. Suppose that you want to permit four consecutive failed connection attempts before the server starts delaying its responses, and to increase the delay by 1500 milliseconds for each additional failure after that. To set the relevant variables at server startup, put these lines in your `my.cnf` file:

```
[mysqld]
plugin-load-add=connection_control.so
connection_control_failed_connections_threshold=4
connection_control_min_connection_delay=1500
```

To set the variables at runtime, use these statements:

```
SET PERSIST connection_control_failed_connections_threshold = 4;
SET PERSIST connection_control_min_connection_delay = 1500;
```

`SET PERSIST` sets the value for the running MySQL instance. It also saves the value to be used for subsequent server restarts; see Section 13.7.4.1, “SET Syntax for Variable Assignment”. To change a value only for the running MySQL instance without saving it for subsequent restarts, use the `GLOBAL` keyword rather than `PERSIST`.

The `connection_control_min_connection_delay` and `connection_control_max_connection_delay` system variables have fixed minimum and maximum values of 1000 and 2147483647, respectively. In addition, the permitted range of values of each variable also depends on the current value of the other:

- `connection_control_min_connection_delay` cannot be set greater than the current value of `connection_control_max_connection_delay`.
- `connection_control_max_connection_delay` cannot be set less than the current value of `connection_control_min_connection_delay`.

Thus, to make the changes required for some configurations, you might need to set the variables in a specific order. Suppose that the current minimum and maximum delays are 1000 and 2000, and that you want to set them to 3000 and 5000. You cannot first set `connection_control_min_connection_delay` to 3000 because that is greater than the current `connection_control_max_connection_delay` value of 2000. Instead, set `connection_control_max_connection_delay` to 5000, then set `connection_control_min_connection_delay` to 3000.

Connection Failure Assessment

When the `CONNECTION_CONTROL` plugin is installed, it checks connection attempts and tracks whether they fail or succeed. For this purpose, a failed connection attempt is one for which the client user and

host match a known MySQL account but the provided credentials are incorrect, or do not match any known account.

Failed-connection counting is based on the user/host combination for each connection attempt. Determination of the applicable user name and host name takes proxying into account and occurs as follows:

- If the client user proxies another user, the proxying user's information is used. For example, if `external_user@example.com` proxies `proxy_user@example.com`, connection counting uses the proxying user, `external_user@example.com`, rather than the proxied user, `proxy_user@example.com`. Both `external_user@example.com` and `proxy_user@example.com` must have valid entries in the `mysql.user` system table and a proxy relationship between them must be defined in the `mysql.proxies_priv` system table (see Section 6.3.11, “Proxy Users”).
- If the client user does not proxy another user, but does match a `mysql.user` entry, counting uses the `CURRENT_USER()` value corresponding to that entry. For example, if a user `user1` connecting from a host `host1.example.com` matches a `user1@host1.example.com` entry, counting uses `user1@host1.example.com`. If the user matches a `user1@%.example.com`, `user1@%.com`, or `user1@%` entry instead, counting uses `user1@%.example.com`, `user1@%.com`, or `user1@%`, respectively.

For the cases just described, the connection attempt matches some `mysql.user` entry, and whether the request succeeds or fails depends on whether the client provides the correct authentication credentials. For example, if the client presents an incorrect password, the connection attempt fails.

If the connection attempt matches no `mysql.user` entry, the attempt fails. In this case, no `CURRENT_USER()` value is available and connection-failure counting uses the user name provided by the client and the client host as determined by the server. For example, if a client attempts to connect as user `user2` from host `host2.example.com`, the user name part is available in the client request and the server determines the host information. The user/host combination used for counting is `user2@host2.example.com`.

Note

The server maintains information about which client hosts can possibly connect to the server (essentially the union of host values for `mysql.user` entries). If a client attempts to connect from any other host, the server rejects the attempt at an early stage of connection setup:

```
ERROR 1130 (HY000): Host 'host_name' is not
allowed to connect to this MySQL server
```

Because this type of rejection occurs so early, `CONNECTION_CONTROL` does not see it, and does not count it.

Connection Failure Monitoring

To monitor failed connections, use these information sources:

- The `Connection_control_delay_generated` status variable indicates the number of times the server added a delay to its response to a failed connection attempt. This does not count attempts that occur before reaching the threshold defined by the `connection_control_failed_connections_threshold` system variable.
- The `INFORMATION_SCHEMA CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` table provides information about the current number of consecutive failed connection attempts per client user/host combination. This counts all failed attempts, regardless of whether they were delayed.

Assigning a value to `connection_control_failed_connections_threshold` at runtime resets all accumulated failed-connection counters to zero, which has these visible effects:

- The `Connection_control_delay_generated` status variable is reset to zero.
- The `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` table becomes empty.

6.5.2.2 Connection-Control System and Status Variables

This section describes the system and status variables that the `CONNECTION_CONTROL` plugin provides to enable its operation to be configured and monitored.

If the `CONNECTION_CONTROL` plugin is installed, it exposes these system variables:

- `connection_control_failed_connections_threshold`

Introduced	8.0.1	
Command-Line Format	`--connection-control-failed-connections-threshold=#`	
System Variable	**Name**	`connection_control_failed_connections_threshold`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`3`
	Min Value	`0`
	Max Value	`2147483647`

The number of consecutive failed connection attempts permitted to clients before the server adds a delay for subsequent connection attempts:

- If the variable has a nonzero value *`N`*, the server adds a delay beginning with consecutive failed attempt *`N`*+1. If a client has reached the point where connection responses are delayed, the delay also occurs for the next subsequent successful connection.
- Setting this variable to zero disables failed-connection counting. In this case, the server never adds delays.

For information about how `connection_control_failed_connections_threshold` interacts with other connection-control system and status variables, see Section 6.5.2.1, “Connection-Control Plugin Installation”.

- `connection_control_max_connection_delay`

Introduced	8.0.1	
Command-Line Format	`--connection-control-max-connection-delay=#`	
System Variable	**Name**	`connection_control_max_connection_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`2147483647`
	Min Value	`1000`

	Max Value	2147483647

The maximum delay in milliseconds for server response to failed connection attempts, if `connection_control_failed_connections_threshold` is greater than zero.

For information about how `connection_control_max_connection_delay` interacts with other connection-control system and status variables, see Section 6.5.2.1, “Connection-Control Plugin Installation”.

- `connection_control_min_connection_delay`

Introduced	8.0.1	
Command-Line Format	`--connection-control-min-connection-delay=#`	
System Variable	**Name**	`connection_control_min_connection_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	1000
	Min Value	1000
	Max Value	2147483647

The minimum delay in milliseconds for server response to failed connection attempts, if `connection_control_failed_connections_threshold` is greater than zero. This is also the amount by which the server increases the delay for additional successive failures once it begins delaying.

For information about how `connection_control_min_connection_delay` interacts with other connection-control system and status variables, see Section 6.5.2.1, “Connection-Control Plugin Installation”.

If the `CONNECTION_CONTROL` plugin is installed, it exposes this status variable:

- `Connection_control_delay_generated`

  The number of times the server added a delay to its response to a failed connection attempt. This does not count attempts that occur before reaching the threshold defined by the `connection_control_failed_connections_threshold` system variable.

  This variable provides a simple counter. For more detailed connection-control monitoring information, examine the `INFORMATION_SCHEMA CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` table; see Section 24.32.1, “The INFORMATION_SCHEMA CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS Table”.

  Assigning a value to `connection_control_failed_connections_threshold` at runtime resets `Connection_control_delay_generated` to zero.

6.5.3 The Password Validation Plugin

The `validate_password` plugin serves to test passwords and improve security. The plugin exposes a set of system variables that enable you to define password policy.

This plugin implements two capabilities:

- In statements that assign a password supplied as a cleartext value, the plugin checks the password against the current password policy and rejects it if it is weak (the statement returns an `ER_NOT_VALID_PASSWORD` error). This affects the `ALTER USER`, `CREATE USER`, `GRANT`, and `SET PASSWORD` statements. Passwords given as arguments to the `PASSWORD()` function are checked as well.
- The `VALIDATE_PASSWORD_STRENGTH()` SQL function assesses the strength of potential passwords. The function takes a password argument and returns an integer from 0 (weak) to 100 (strong).

For example, the cleartext password in the following statement is checked. Under the default password policy, which requires passwords to be at least 8 characters long, the password is weak and the statement produces an error:

```
mysql> ALTER USER USER() IDENTIFIED BY 'abc';
ERROR 1819 (HY000): Your password does not satisfy the current
policy requirements
```

Passwords specified as hashed values are not checked because the original password value is not available:

```
mysql> ALTER USER 'jeffrey'@'localhost'
       IDENTIFIED WITH mysql_native_password
       AS '*0D3CED9BEC10A777AEC23CCC353A8C08A633045E';
Query OK, 0 rows affected (0.01 sec)
```

System variables having names of the form `validate_password_xxx` represent the parameters that control password policy. To configure password checking, modify these variables; see Section 6.5.3.2, “Password Validation Plugin Options and Variables”.

If the `validate_password` plugin is not installed, the `validate_password_xxx` system variables are not available, passwords in statements are not checked, and the `VALIDATE_PASSWORD_STRENGTH()` function always returns 0. For example, without the plugin installed, accounts can be assigned passwords shorter than 8 characters.

Assuming that the `validate_password` plugin is installed, it implements three levels of password checking: `LOW`, `MEDIUM`, and `STRONG`. The default is `MEDIUM`; to change this, modify the value of `validate_password_policy`. The policies implement increasingly strict password tests. The following descriptions refer to default parameter values, which can be modified by changing the appropriate system variables.

- `LOW` policy tests password length only. Passwords must be at least 8 characters long.
- `MEDIUM` policy adds the conditions that passwords must contain at least 1 numeric character, 1 lowercase character, 1 uppercase character, and 1 special (nonalphanumeric) character.
- `STRONG` policy adds the condition that password substrings of length 4 or longer must not match words in the dictionary file, if one has been specified.

In addition, the `validate_password` plugin supports the capability of rejecting passwords that match the user name part of the effective user account for the current session, either forward or in reverse. To enable control over this capability, the plugin exposes a `validate_password_check_user_name` system variable. By default, this variable is enabled.

6.5.3.1 Password Validation Plugin Installation

This section describes how to install the `validate_password` password-validation plugin. For general information about installing plugins, see Section 6.6.2, “Installing and Uninstalling Plugins”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base name is `validate_password`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To load the plugin at server startup, use the `--plugin-load-add` option to name the library file that contains it. With this plugin-loading method, the option must be given each time the server starts. For example, put these lines in your `my.cnf` file (adjust the `.so` suffix for your platform as necessary):

```
[mysqld]
plugin-load-add=validate_password.so
```

Alternatively, to register the plugin at runtime, use this statement (adjust the extension as necessary):

```
INSTALL PLUGIN validate_password SONAME 'validate_password.so';
```

`INSTALL PLUGIN` loads the plugin, and also registers it in the `mysql.plugins` table to cause the plugin to be loaded for each subsequent normal server startup.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE 'validate%';
+-------------------+---------------+
| PLUGIN_NAME       | PLUGIN_STATUS |
+-------------------+---------------+
| validate_password | ACTIVE        |
+-------------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

If the plugin has been previously registered with `INSTALL PLUGIN` or is loaded with `--plugin-load-add`, you can use the `--validate-password` option at server startup to control plugin activation. For example, to load the plugin at startup and prevent it from being removed at runtime, use these options:

```
[mysqld]
plugin-load-add=validate_password.so
validate-password=FORCE_PLUS_PERMANENT
```

If it is desired to prevent the server from running without the password-validation plugin, use `--validate-password` with a value of `FORCE` or `FORCE_PLUS_PERMANENT` to force server startup to fail if the plugin does not initialize successfully.

6.5.3.2 Password Validation Plugin Options and Variables

To control the activation of the `validate_password` plugin, use this option:

- `--validate-password[=value]`

Command-Line Format	`--validate-password[=value]`	
Permitted Values	**Type**	enumeration
	Default	`ON`
	Valid Values	`ON`
		`OFF`
		`FORCE`
		`FORCE_PLUS_PERMANENT`

This option controls how the server loads the `validate_password` plugin at startup. The value should be one of those available for plugin-loading options, as described in Section 5.6.2, "Installing and Uninstalling Plugins". For example, `--validate-password=FORCE_PLUS_PERMANENT` tells the server to load the plugin at startup and prevents it from being removed while the server is running.

This option is available only if the `validate_password` plugin has been previously registered with `INSTALL PLUGIN` or is loaded with `--plugin-load-add`. See Section 6.5.3.1, "Password Validation Plugin Installation".

If the `validate_password` plugin is enabled, it exposes several system variables representing the parameters that control password checking:

```
mysql> SHOW VARIABLES LIKE 'validate_password%';
+--------------------------------------+--------+
| Variable_name                        | Value  |
+--------------------------------------+--------+
| validate_password_check_user_name    | ON     |
| validate_password_dictionary_file    |        |
| validate_password_length             | 8      |
| validate_password_mixed_case_count   | 1      |
| validate_password_number_count       | 1      |
| validate_password_policy             | MEDIUM |
| validate_password_special_char_count | 1      |
+--------------------------------------+--------+
```

To change how passwords are checked, you can set these system variables at server startup or at runtime. The following list describes the meaning of each variable.

- `validate_password_check_user_name`

Command-Line Format		`--validate-password-check-user-name`
System Variable	Name	`validate_password_check_user_name`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

Whether passwords are compared to the user name part of the effective user account for the current session and rejected if they match. By default, `validate_password_check_user_name` is enabled. This variable controls user name matching independent of the value of `validate_password_policy`.

When `validate_password_check_user_name` is enabled, it has these effects:

- Checking occurs in all contexts for which the `validate_password` plugin is invoked, which includes statements such as `ALTER USER` and `SET PASSWORD`, and invocation of functions such as `PASSWORD()` and `VALIDATE_PASSWORD_STRENGTH()`.
- If a password is the same as the user name or its reverse, a match occurs and the password is rejected.
- If a password matches the user name, `VALIDATE_PASSWORD_STRENGTH()` returns 0 regardless of how other `validate_password` system variables are set.
- The user names used for comparison are taken from the values of the `USER()` and `CURRENT_USER()` functions for the current session. (An implication is that a user who has the

SUPER privilege can execute a statement to set another user's password to that user name, and cannot set that user's password to the name of the user executing the statement.)

- Only the user name part of the USER() and CURRENT_USER() function values is used, not the host name part. If a user name is empty, no comparison is done.
- User name matching is case sensitive. The password and user name values are compared as binary strings on a byte-by-byte basis.

- validate_password_dictionary_file

System Variable	Name	validate_password_dictionary_file
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	file name

The path name of the dictionary file used by the validate_password plugin for checking passwords. This variable is unavailable unless that plugin is installed.

By default, this variable has an empty value and dictionary checks are not performed. To enable dictionary checks, you must set this variable to a nonempty value. If the file is named as a relative path, it is interpreted relative to the server data directory. Its contents should be lowercase, one word per line. Contents are treated as having a character set of utf8. The maximum permitted file size is 1MB.

For the dictionary file to be used during password checking, the password policy must be set to 2 (STRONG); see the description of the validate_password_policy system variable. Assuming that is true, each substring of the password of length 4 up to 100 is compared to the words in the dictionary file. Any match causes the password to be rejected. Comparisons are not case sensitive.

For VALIDATE_PASSWORD_STRENGTH(), the password is checked against all policies, including STRONG, so the strength assessment includes the dictionary check regardless of the validate_password_policy value.

validate_password_dictionary_file can be set at runtime and assigning a value causes the named file to be read without a restart.

- validate_password_length

System Variable	Name	validate_password_length
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	8
	Min Value	0

The minimum number of characters that passwords checked by the validate_password plugin must have. This variable is unavailable unless that plugin is installed.

The validate_password_length minimum value is a function of several other related system variables. The server will not set the value less than the value of this expression:

```
validate_password_number_count
+ validate_password_special_char_count
+ (2 * validate_password_mixed_case_count)
```

If the `validate_password` plugin adjusts the value of `validate_password_length` due to the preceding constraint, it writes a message to the error log.

- `validate_password_mixed_case_count`

System Variable	Name	validate_password_mixed_case_count
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	1
	Min Value	0

The minimum number of lowercase and uppercase characters that passwords checked by the `validate_password` plugin must have if the password policy is `MEDIUM` or stronger. For a given value, the password must have that many lowercase characters, and that many uppercase characters. This variable is unavailable unless that plugin is installed.

- `validate_password_number_count`

System Variable	Name	validate_password_number_count
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	1
	Min Value	0

The minimum number of numeric (digit) characters that passwords checked by the `validate_password` plugin must have if the password policy is `MEDIUM` or stronger. This variable is unavailable unless that plugin is installed.

- `validate_password_policy`

System Variable	Name	validate_password_policy
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	1
	Valid Values	0
		1

		2

The password policy enforced by the `validate_password` plugin. This variable is unavailable unless that plugin is installed.

`validate_password_policy` affects how the plugin uses its other policy-setting system variables, except for checking passwords against user names, which is controlled independently by `validate_password_check_user_name`.

The `validate_password_policy` value can be specified using numeric values 0, 1, 2, or the corresponding symbolic values `LOW`, `MEDIUM`, `STRONG`. The following table describes the tests performed for each policy. For the length test, the required length is the value of the `validate_password_length` system variable. Similarly, the required values for the other tests are given by other `validate_password_xxx` variables.

Policy	Tests Performed
`0` or `LOW`	Length
`1` or `MEDIUM`	Length; numeric, lowercase/uppercase, and special characters
`2` or `STRONG`	Length; numeric, lowercase/uppercase, and special characters; dictionary file

- `validate_password_special_char_count`

System Variable	Name	`validate_password_special_char_count`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1`
	Min Value	`0`

The minimum number of nonalphanumeric characters that passwords checked by the `validate_password` plugin must have if the password policy is `MEDIUM` or stronger. This variable is unavailable unless that plugin is installed.

If the `validate_password` plugin is enabled, it exposes status variables that provide operational information:

```
mysql> SHOW STATUS LIKE 'validate_password%';
+----------------------------------------------+---------------------+
| Variable_name                                | Value               |
+----------------------------------------------+---------------------+
| validate_password_dictionary_file_last_parsed | 2015-06-29 11:08:51 |
| validate_password_dictionary_file_words_count | 1902                |
+----------------------------------------------+---------------------+
```

The following list describes the meaning of each status variable.

- `validate_password_dictionary_file_last_parsed`

  When the dictionary file was last parsed.

- `validate_password_dictionary_file_words_count`

  The number of words read from the dictionary file.

6.5.4 The MySQL Keyring

MySQL Server supports a keyring service that enables internal server components and plugins to securely store sensitive information for later retrieval. The implementation is plugin-based:

- The `keyring_file` plugin stores keyring data in a file local to the server host. This plugin is available in all MySQL distributions, Community Edition and Enterprise Edition included.

Warning

The `keyring_file` plugin for encryption key management is not intended as a regulatory compliance solution. Security standards such as PCI, FIPS, and others require use of key management systems to secure, manage, and protect encryption keys in key vaults or hardware security modules (HSMs).

- An SQL interface for keyring key management is implemented as a set of user-defined functions (UDFs).

The `InnoDB` storage engine uses a keyring plugin store its key for tablespace encryption. `InnoDB` can use any supported keyring plugin.

For keyring installation instructions, see Section 6.5.4.1, “Keyring Plugin Installation”.

For information about using the keyring UDFs, see Section 6.5.4.4, “General-Purpose Keyring Key-Management Functions”.

Keyring plugins and UDFs access a keyring service that provides the interface for server components to the keyring. For information about writing keyring plugins or accessing the keyring plugin service, see Section 28.2.4.12, “Writing Keyring Plugins”, and Section 28.3.2, “The Keyring Service”.

6.5.4.1 Keyring Plugin Installation

Keyring service consumers require a keyring plugin to be installed. MySQL provides these plugin choices:

- `keyring_file`: A plugin that stores keyring data in a file local to the server host. Available in all MySQL distributions.

This section describes how to install the keyring plugin of your choosing. For general information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

Installation for each keyring plugin is similar. The following instructions use `keyring_file`. Users of a different keyring plugin can substitute its name for `keyring_file`.

The `keyring_file` plugin library file base name is `keyring_file`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

Note

Only one keyring plugin should be enabled at a time. Enabling multiple keyring plugins is unsupported and results may not be as anticipated.

The keyring plugin must be loaded early during the server startup sequence so that server components can access it as necessary during their own initialization. For example, the `InnoDB` storage engine uses the keyring for tablespace encryption, so the keyring plugin must be loaded and available prior to `InnoDB` initialization.

To load the plugin, use the `--early-plugin-load` option to name the plugin library file that contains it. For example, on platforms where the plugin library file suffix is `.so`, use these lines in the server `my.cnf` file (adjust the `.so` suffix for your platform as necessary):

```
[mysqld]
early-plugin-load=keyring_file.so
```

Before starting the server, check the notes for your chosen keyring plugin to see whether it permits or requires additional configuration:

- For `keyring_file`: Section 6.5.4.2, “Using the keyring_file File-Based Plugin”.

After performing any plugin-specific configuration, verify plugin installation. With the MySQL server running, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE 'keyring%';
+--------------+---------------+
| PLUGIN_NAME  | PLUGIN_STATUS |
+--------------+---------------+
| keyring_file | ACTIVE        |
+--------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

If no keyring plugin is available when a server component tries to access the keyring service, the service cannot be used by that component. As a result, the component may fail to initialize or may initialize with limited functionality. For example, if `InnoDB` finds that there are encrypted tablespaces when it initializes, it attempts to access the keyring. If the keyring is unavailable, `InnoDB` can access only unencrypted tablespaces. To ensure that `InnoDB` can access encrypted tablespaces as well, use `--early-plugin-load` to load the keyring plugin.

Plugins can be loaded by other methods, such as the `--plugin-load` or `--plugin-load-add` option or the `INSTALL PLUGIN` statement. However, keyring plugins loaded using those methods may be available too late in the server startup sequence for certain server components, such as `InnoDB`:

- Plugin loading using `--plugin-load` or `--plugin-load-add` occurs after `InnoDB` initialization.
- Plugins installed using `INSTALL PLUGIN` are registered in the `mysql.plugin` system table and loaded automatically for subsequent server restarts. However, because `mysql.plugin` is an `InnoDB` table, any plugins named in it can be loaded during startup only after `InnoDB` initialization.

6.5.4.2 Using the keyring_file File-Based Plugin

The `keyring_file` plugin is a keyring plugin that stores keyring data in a file local to the server host.

Warning

The `keyring_file` plugin for encryption key management is not intended as a regulatory compliance solution. Security standards such as PCI, FIPS, and others require use of key management systems to secure, manage, and protect encryption keys in key vaults or hardware security modules (HSMs).

To install the `keyring_file` plugin, use the general keyring installation instructions found in Section 6.5.4.1, “Keyring Plugin Installation”, together with the configuration information specific to `keyring_file` found here.

`keyring_file` must be loaded at each server startup using the `--early-plugin-load` option. The `keyring_file_data` system variable optionally configures the location of the file used by the

`keyring_file` plugin for data storage. The default value is platform specific. To configure the file location explicitly, set the variable value at startup. For example, use these lines in the server `my.cnf` file (adjust the `.so` suffix and file location for your platform as necessary):

```
[mysqld]
early-plugin-load=keyring_file.so
keyring_file_data=/usr/local/mysql/mysql-keyring/keyring
```

Keyring operations are transactional: The `keyring_file` plugin uses a backup file during write operations to ensure that it can roll back to the original file if an operation fails. The backup file has the same name as the value of the `keyring_file_data` system variable with an extension of `.backup`.

For additional information about `keyring_file_data`, see Section 6.5.4.6, “Keyring System Variables”.

To ensure that keys are flushed only when the correct keyring storage file exists, `keyring_file` stores a SHA-256 checksum of the keyring in the file. Before updating the file, the plugin verifies that it contains the expected checksum.

The `keyring_file` plugin supports the functions that comprise the standard keyring service interface. Keyring operations performed by these functions are accessible at two levels:

- SQL interface: In SQL statements, call the user-defined functions (UDFs) described in Section 6.5.4.4, “General-Purpose Keyring Key-Management Functions”.
- C interface: In C-language code, call the keyring service functions described in Section 28.3.2, “The Keyring Service”.

Example (using UDFs):

```
SELECT keyring_key_generate('MyKey', 'AES', 32);
SELECT keyring_key_remove('MyKey');
```

The key types permitted by `keyring_file` are described in Section 6.5.4.3, “Supported Keyring Key Types”.

6.5.4.3 Supported Keyring Key Types

MySQL Keyring supports generating keys of different types (encryption algorithms) and lengths. The available key types depend on which keyring plugin is installed. A given plugin may also impose constraints on key lengths per key type.

Table 6.16, “Keyring Plugin Key Types” summarizes the permitted key types per keyring plugin. Lengths are in bytes. For a key generated using one of the keyring user-defined functions (UDFs) described in Section 6.5.4.4, “General-Purpose Keyring Key-Management Functions”, the length can be no more than 2,048 bytes, due to limitations of the UDF interface.

Table 6.16 Keyring Plugin Key Types

Plugin Name	Permitted Key Type	Permitted Key Lengths for Key Type
`keyring_file`	AES	No special restrictions
	DSA	No special restrictions
	RSA	No special restrictions

6.5.4.4 General-Purpose Keyring Key-Management Functions

MySQL Server supports a keyring service that enables internal server components and plugins to securely store sensitive information for later retrieval.

MySQL Server also includes an SQL interface for keyring key management, implemented as a set of general-purpose user-defined functions (UDFs) that access the functions provided by the internal keyring service. The keyring UDFs are contained in a plugin library file, which also contains a `keyring_udf` plugin that must be enabled prior to UDF invocation. For these UDFs to be used, a keyring plugin such as `keyring_file` must be enabled.

The UDFs described here are general purpose and intended for use with any keyring plugin. A given keyring plugin might have UDFs of its own that are intended for use only with that plugin; see Section 6.5.4.5, “Plugin-Specific Keyring Key-Management Functions”.

The following sections provide installation instructions for the keyring UDFs and demonstrate how to use them. For information about the keyring service functions invoked by the UDFs, see Section 28.3.2, “The Keyring Service”. For general keyring information, see Section 6.5.4, “The MySQL Keyring”.

Installing or Uninstalling General-Purpose Keyring Functions

This section describes how to install or uninstall the keyring user-defined functions (UDFs), which are implemented in a plugin library file that also contains a `keyring_udf` plugin. For general information about installing or uninstalling plugins and UDFs, see Section 5.6.2, “Installing and Uninstalling Plugins”, and Section 28.4.2.5, “UDF Compiling and Installing”.

The keyring UDFs enable keyring key management operations, but the `keyring_udf` plugin must also be installed because the UDFs will not work correctly without it. Attempts to use the UDFs without the `keyring_udf` plugin result in an error.

To be usable by the server, the plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base name is `keyring_udf`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To install the `keyring_udf` plugin and the UDFs, use the `INSTALL PLUGIN` and `CREATE FUNCTION` statements (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN keyring_udf SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_generate RETURNS INTEGER SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_fetch RETURNS STRING SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_length_fetch RETURNS INTEGER SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_type_fetch RETURNS STRING SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_store RETURNS INTEGER SONAME 'keyring_udf.so';
CREATE FUNCTION keyring_key_remove RETURNS INTEGER SONAME 'keyring_udf.so';
```

If the plugin and the UDFs are used on a master replication server, install them on all slave servers as well to avoid replication problems.

Once installed as just described, the `keyring_udf` plugin and the UDFs remain installed until uninstalled. To remove them, use the `UNINSTALL PLUGIN` and `DROP FUNCTION` statements:

```
UNINSTALL PLUGIN keyring_udf;
DROP FUNCTION keyring_key_generate;
DROP FUNCTION keyring_key_fetch;
DROP FUNCTION keyring_key_length_fetch;
DROP FUNCTION keyring_key_type_fetch;
DROP FUNCTION keyring_key_store;
DROP FUNCTION keyring_key_remove;
```

Using General-Purpose Keyring Functions

Before using the keyring user-defined functions (UDFs), install them according to the instructions provided in Installing or Uninstalling General-Purpose Keyring Functions.

The keyring UDFs are subject to these constraints:

- To use any keyring UDF, the `keyring_udf` plugin must be enabled. Otherwise, an error occurs:

```
ERROR 1123 (HY000): Can't initialize function 'keyring_key_generate';
This function requires keyring_udf plugin which is not installed.
Please install
```

  To install the `keyring_udf` plugin, see Installing or Uninstalling General-Purpose Keyring Functions.

- The keyring UDFs invoke keyring service functions (see Section 28.3.2, “The Keyring Service”). The service functions in turn use whatever keyring plugin is installed (for example, `keyring_file`). Therefore, to use any keyring UDF, some underlying keyring plugin must be enabled. Otherwise, an error occurs:

```
ERROR 3188 (HY000): Function 'keyring_key_generate' failed because
underlying keyring service returned an error. Please check if a
keyring plugin is installed and that provided arguments are valid
for the keyring you are using.
```

  To install a keyring plugin, see Section 6.5.4.1, “Keyring Plugin Installation”.

- To use any keyring UDF, a user must possess the global `EXECUTE` privilege. Otherwise, an error occurs:

```
ERROR 1123 (HY000): Can't initialize function 'keyring_key_generate';
The user is not privileged to execute this function. User needs to
have EXECUTE
```

  To grant the global `EXECUTE` privilege to a user, use this statement:

```
GRANT EXECUTE ON *.* TO user;
```

  Alternatively, should you prefer to avoid granting the global `EXECUTE` privilege while still permitting users to access specific key-management operations, “wrapper” stored programs can be defined (a technique described later in this section).

- A key stored in the keyring by a given user can be manipulated later only by the same user. That is, the value of the `CURRENT_USER()` function at the time of key manipulation must have the same value as when the key was stored in the keyring. (This constraint rules out the use of the keyring UDFs for manipulation of instance-wide keys, such as those created by `InnoDB` to support tablespace encryption.)

  To enable multiple users to perform operations on the same key, “wrapper” stored programs can be defined (a technique described later in this section).

- Keyring UDFs support the key types and lengths supported by the underlying keyring plugin, with the additional constraint that keys cannot be longer than 2,048 bytes (16,384 bits), due to limitations of the UDF interface. See Section 6.5.4.3, “Supported Keyring Key Types”.

To create a new random key and store it in the keyring, call `keyring_key_generate()`, passing to it an ID for the key, along with the key type (encryption method) and its length in bytes. The following call creates a 2,048-bit DSA-encrypted key named `MyKey`:

```
mysql> SELECT keyring_key_generate('MyKey', 'DSA', 256);
+-----------------------------------------+
| keyring_key_generate('MyKey', 'DSA', 256) |
+-----------------------------------------+
|                                       1 |
+-----------------------------------------+
```

A return value of 1 indicates success. If the key cannot be created, the return value is `NULL` and an error occurs. To be able to check the return type regardless of whether an error occurs, use `SELECT ... INTO @var_name` and test the variable value:

```
mysql> SELECT keyring_key_generate('', '', -1) INTO @x;
ERROR 3188 (HY000): Function 'keyring_key_generate' failed because
underlying keyring service returned an error. Please check if a
keyring plugin is installed and that provided arguments are valid
for the keyring you are using.
mysql> SELECT @x;
+------+
| @x   |
+------+
| NULL |
+------+
mysql> SELECT keyring_key_generate('x', 'AES', 16) INTO @x;
mysql> SELECT @x;
+------+
| @x   |
+------+
|    1 |
+------+
```

This technique also applies to other keyring UDFs that for failure return a value and an error.

The ID passed to `keyring_key_generate()` provides a means by which to refer to the key in subsequent UDF calls. For example, use the key ID to retrieve its type as a string or its length in bytes as an integer:

```
mysql> SELECT keyring_key_type_fetch('MyKey');
+---------------------------------+
| keyring_key_type_fetch('MyKey') |
+---------------------------------+
| DSA                             |
+---------------------------------+
mysql> SELECT keyring_key_length_fetch('MyKey');
+----------------------------------+
| keyring_key_length_fetch('MyKey') |
+----------------------------------+
|                              256 |
+----------------------------------+
```

To retrieve a key value, pass the key ID to `keyring_key_fetch()`. The following example uses `HEX()` to display the key value because it may contain nonprintable characters. The example also uses a short key for brevity, but be aware that longer keys provide better security:

```
mysql> SELECT keyring_key_generate('MyShortKey', 'DSA', 8);
+---------------------------------------------+
| keyring_key_generate('MyShortKey', 'DSA', 8) |
+---------------------------------------------+
|                                           1 |
+---------------------------------------------+
mysql> SELECT HEX(keyring_key_fetch('MyShortKey'));
+-------------------------------------+
| HEX(keyring_key_fetch('MyShortKey')) |
+-------------------------------------+
| 1DB3B0FC3328A24C                    |
+-------------------------------------+
```

Keyring UDFs treat key IDs, types, and values as binary strings, so comparisons are case sensitive. For example, IDs of `MyKey` and `mykey` refer to different keys.

To remove a key, pass the key ID to `keyring_key_remove()`:

```
mysql> SELECT keyring_key_remove('MyKey');
```

```
+----------------------------+
| keyring_key_remove('MyKey') |
+----------------------------+
|                          1 |
+----------------------------+
```

To obfuscate and store a key that you provide, pass the key ID, type, and value to `keyring_key_store()`:

```
mysql> SELECT keyring_key_store('AES_key', 'AES', 'Secret string');
+-----------------------------------------------------+
| keyring_key_store('AES_key', 'AES', 'Secret string') |
+-----------------------------------------------------+
|                                                   1 |
+-----------------------------------------------------+
```

As indicated previously, a user must have the global `EXECUTE` privilege to call keyring UDFs, and the user who stores a key in the keyring initially must be the same user who performs subsequent operations on the key later, as determined from the `CURRENT_USER()` value in effect for each UDF call. To permit key operations to users who do not have the global `EXECUTE` privilege or who may not be the key "owner," use this technique:

1. Define "wrapper" stored programs that encapsulate the required key operations and have a `DEFINER` value equal to the key owner.
2. Grant the `EXECUTE` privilege for specific stored programs to the individual users who should be able to invoke them.
3. If the operations implemented by the wrapper stored programs do not include key creation, create any necessary keys in advance, using the account named as the `DEFINER` in the stored program definitions.

This technique enables keys to be shared among users and provides to DBAs more fine-grained control over who can do what with keys, without having to grant global privileges.

The following example shows how to set up a shared key named `SharedKey` that is owned by the DBA, and a `get_shared_key()` stored function that provides access to the current key value. The value can be retrieved by any user with the `EXECUTE` privilege for that function, which is created in the `key_schema` schema.

From a MySQL administrative account (`'root'@'localhost'` in this example), create the administrative schema and the stored function to access the key:

```
mysql> CREATE SCHEMA key_schema;

mysql> CREATE DEFINER = 'root'@'localhost'
       FUNCTION key_schema.get_shared_key()
       RETURNS BLOB READS SQL DATA
       RETURN keyring_key_fetch('SharedKey');
```

From the administrative account, ensure that the shared key exists:

```
mysql> SELECT keyring_key_generate('SharedKey', 'DSA', 8);
+-------------------------------------------+
| keyring_key_generate('SharedKey', 'DSA', 8) |
+-------------------------------------------+
|                                         1 |
+-------------------------------------------+
```

From the administrative account, create an ordinary user account to which key access is to be granted:

```
```

```
mysql> CREATE USER 'key_user'@'localhost'
       IDENTIFIED BY 'key_user_pwd';
```

From the `key_user` account, verify that, without the proper `EXECUTE` privilege, the new account cannot access the shared key:

```
mysql> SELECT HEX(key_schema.get_shared_key());
ERROR 1370 (42000): execute command denied to user 'key_user'@'localhost'
for routine 'key_schema.get_shared_key'
```

From the administrative account, grant `EXECUTE` to `key_user` for the stored function:

```
mysql> GRANT EXECUTE ON FUNCTION key_schema.get_shared_key
       TO 'key_user'@'localhost';
```

From the `key_user` account, verify that the key is now accessible:

```
mysql> SELECT HEX(key_schema.get_shared_key());
+---------------------------------+
| HEX(key_schema.get_shared_key()) |
+---------------------------------+
| 9BAFB9E75CEEB013                |
+---------------------------------+
```

General-Purpose Keyring Function Reference

For each general-purpose keyring user-defined function (UDF), this section describes its purpose, calling sequence, and return value. For information about the conditions under which these UDFs can be invoked, see Using General-Purpose Keyring Functions.

- `keyring_key_fetch()`

  Given a key ID, deobfuscates and returns the key value.

  Syntax:

  ```
  STRING keyring_key_fetch(STRING key_id)
  ```

  Arguments:

  - `key_id`: The key ID as a string.

  Return values:

  Returns the key value for success, `NULL` if the key does not exist, or `NULL` and an error for failure.

  Note

  Keyring values retrieved using `keyring_key_fetch()` are limited to 2,048 bytes, due to limitations of the UDF interface. A keyring value longer than that length can be stored using a keyring service function (see Section 28.3.2, “The Keyring Service”), but if retrieved using `keyring_key_fetch()`, is truncated to 2,048 bytes.

  Example:

  ```
  mysql> SELECT keyring_key_generate('RSA_key', 'RSA', 16);
  +-------------------------------------------+
  | keyring_key_generate('RSA_key', 'RSA', 16) |
  +-------------------------------------------+
  |                                         1 |
  +-------------------------------------------+
  ```

```
mysql> SELECT HEX(keyring_key_fetch('RSA_key'));
+----------------------------------+
| HEX(keyring_key_fetch('RSA_key')) |
+----------------------------------+
| 91C2253B696064D3556984B6630F891A  |
+----------------------------------+
mysql> SELECT keyring_key_type_fetch('RSA_key');
+----------------------------------+
| keyring_key_type_fetch('RSA_key') |
+----------------------------------+
| RSA                              |
+----------------------------------+
mysql> SELECT keyring_key_length_fetch('RSA_key');
+-----------------------------------+
| keyring_key_length_fetch('RSA_key') |
+-----------------------------------+
|                                16 |
+-----------------------------------+
```

The example uses `HEX()` to display the key value because it may contain nonprintable characters. The example also uses a short key for brevity, but be aware that longer keys provide better security.

- `keyring_key_generate()`

  Generates a new random key with a given ID, type, and length, and stores it in the keyring. The type and length values must be consistent with the values supported by the underlying keyring plugin, with the additional constraint that keys cannot be longer than 2,048 bytes (16,384 bits), due to limitations of the UDF interface. For the permitted types per plugin, see Section 28.3.2, “The Keyring Service”.

  Syntax:

  ```
  STRING keyring_key_generate(STRING key_id, STRING key_type, INTEGER key_length)
  ```

  Arguments:

  - `key_id`: The key ID as a string.
  - `key_type`: The key type as a string.
  - `key_length`: The key length in bytes as an integer. The maximum length is 2,048.

  Return values:

  Returns 1 for success, or `NULL` and an error for failure.

  Example:

  ```
  mysql> SELECT keyring_key_generate('RSA_key', 'RSA', 384);
  +--------------------------------------------+
  | keyring_key_generate('RSA_key', 'RSA', 384) |
  +--------------------------------------------+
  |                                          1 |
  +--------------------------------------------+
  ```

- `keyring_key_length_fetch()`

  Given a key ID, returns the key length.

  Syntax:

  ```
  INTEGER keyring_key_length_fetch(STRING key_id)
  ```

  Arguments:

- `key_id`: The key ID as a string.

Return values:

Returns the key length in bytes as an integer for success, `NULL` if the key does not exist, or `NULL` and an error for failure.

Example:

See the description of `keyring_key_fetch()`.

- `keyring_key_remove()`

  Removes the key with a given ID from the keyring.

  Syntax:

```
INTEGER keyring_key_remove(STRING key_id)
```

  Arguments:

  - `key_id`: The key ID as a string.

  Return values:

  Returns 1 for success, or `NULL` for failure.

  Example:

```
mysql> SELECT keyring_key_remove('AES_key');
+------------------------------+
| keyring_key_remove('AES_key') |
+------------------------------+
|                            1 |
+------------------------------+
```

- `keyring_key_store()`

  Obfuscates and stores a key in the keyring.

  Syntax:

```
INTEGER keyring_key_store(STRING key_id, STRING key_type, STRING key)
```

  Arguments:

  - `key_id`: The key ID as a string.
  - `key_type`: The key type as a string.
  - `key`: The key value as a string.

  Return values:

  Returns 1 for success, or `NULL` and an error for failure.

  Example:

```
mysql> SELECT keyring_key_store('new key', 'DSA', 'My key value');
+-------------------------------------------------+
| keyring_key_store('new key', 'DSA', 'My key value') |
```

```
+--------------------------------------------------+
|                                                1 |
+--------------------------------------------------+
```

- `keyring_key_type_fetch()`

  Given a key ID, returns the key type.

  Syntax:

  ```
  STRING keyring_key_type_fetch(STRING key_id)
  ```

  Arguments:

  - `key_id`: The key ID as a string.

  Return values:

  Returns the key type as a string for success, `NULL` if the key does not exist, or `NULL` and an error for failure.

  Example:

  See the description of `keyring_key_fetch()`.

6.5.4.5 Plugin-Specific Keyring Key-Management Functions

This section describes user-defined functions (UDFs) that are specific to individual keyring plugins. Currently, no keyring plugin in MySQL 8.0 has plugin-specific UDFs. For information about general-purpose keyring UDFs, see Section 6.5.4.4, “General-Purpose Keyring Key-Management Functions”.

6.5.4.6 Keyring System Variables

MySQL Keyring plugins support the following system variables. Use them to configure keyring plugin operation. These variables are unavailable unless the appropriate keyring plugin is installed (see Section 6.5.4.1, “Keyring Plugin Installation”).

- `keyring_file_data`

Command-Line Format	`--keyring-file-data=file_name`	
System Variable	**Name**	`keyring_file_data`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	file name
	Default	`platform specific`

The path name of the data file used for secure data storage by the `keyring_file` plugin. This variable is unavailable unless that plugin is installed. The file location should be in a directory considered for use only by the `keyring_file` plugin. For example, do not locate the file under the data directory.

Keyring operations are transactional: The `keyring_file` plugin uses a backup file during write operations to ensure that it can roll back to the original file if an operation fails. The backup file has the same name as the value of the `keyring_file_data` system variable with an extension of `.backup`.

Do not use the same `keyring_file` data file for multiple MySQL instances. Each instance should have its own unique data file.

The default file name is `keyring`, located in a directory that is platform specific and depends on the value of the `INSTALL_LAYOUT` `CMake` option, as shown in the following table. To specify the default directory for the file explicitly if you are building from source, use the `INSTALL_MYSQLKEYRINGDIR` `CMake` option.

`INSTALL_LAYOUT` Value	**Default `keyring_file_data` Value**
`DEB`, `RPM`, `SLES`, `SVR4`	`/var/lib/mysql-keyring/keyring`
Otherwise	`keyring/keyring` under the `CMAKE_INSTALL_PREFIX` value

At plugin startup, if the value assigned to `keyring_file_data` specifies a file that does not exist, the `keyring_file` plugin attempts to create it (as well as its parent directory, if necessary).

If you create the directory manually, it should have a restrictive mode and be accessible only to the account used to run the MySQL server. For example, on Unix and Unix-like systems, to use `/usr/local/mysql/mysql-keyring/keyring`, the following commands (executed as `root`) create the directory and set its mode and ownership:

```
cd /usr/local/mysql
mkdir mysql-keyring
chmod 750 mysql-keyring
chown mysql mysql-keyring
chgrp mysql mysql-keyring
```

If the `keyring_file` plugin cannot create or access the file, it writes an error message to the error log. If an attempted runtime assignment to `keyring_file_data` results in an error, the variable value remains unchanged.

Important

Once the `keyring_file` plugin has created the `keyring_file` plugin data file and started to use it, it is important not to remove the file. For example, `InnoDB` uses the file to store the master key used to decrypt the data in tables that use `InnoDB` tablespace encryption; see Section 15.7.10, “InnoDB Tablespace Encryption”. Loss of the file will cause data in such tables to become inaccessible. (It is permissible to rename or move the file, as long as you change the value of `keyring_file_data` to match.) It is recommended that you create a separate backup of the `keyring` file immediately after you create the first encrypted table and before and after master key rotation.

Chapter 7 Backup and Recovery

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It is important to back up your databases so that you can recover your data and be up and running again in case problems occur, such as system crashes, hardware failures, or users deleting data by mistake. Backups are also essential as a safeguard before upgrading a MySQL installation, and they can be used to transfer a MySQL installation to another system or to set up replication slave servers.

MySQL offers a variety of backup strategies from which you can choose the methods that best suit the requirements for your installation. This chapter discusses several backup and recovery topics with which you should be familiar:

- Types of backups: Logical versus physical, full versus incremental, and so forth.
- Methods for creating backups.
- Recovery methods, including point-in-time recovery.
- Backup scheduling, compression, and encryption.
- Table maintenance, to enable recovery of corrupt tables.

Additional Resources

Resources related to backup or to maintaining data availability include the following:

- Customers of MySQL Enterprise Edition can use the MySQL Enterprise Backup product for backups. For an overview of the MySQL Enterprise Backup product, see Section 29.2, “MySQL Enterprise Backup Overview”.
- A forum dedicated to backup issues is available at http://forums.mysql.com/list.php?28.
- Details for `mysqldump` can be found in Chapter 4, *MySQL Programs*.
- The syntax of the SQL statements described here is given in Chapter 13, *SQL Statement Syntax*.

- For additional information about `InnoDB` backup procedures, see Section 15.17.1, “InnoDB Backup”.
- Replication enables you to maintain identical data on multiple servers. This has several benefits, such as enabling client query load to be distributed over servers, availability of data even if a given server is taken offline or fails, and the ability to make backups with no impact on the master by using a slave server. See Chapter 18, *Replication*.
- MySQL InnoDB cluster is a collection of products that work together to provide a high availability solution. A group of MySQL servers can be configured to create a cluster using MySQL Shell. The cluster of servers has a single master, called the primary, which acts as the read-write master. Multiple secondary servers are replicas of the master. A minimum of three servers are required to create a high availability cluster. A client application is connected to the primary via MySQL Router. If the primary fails, a secondary is automatically promoted to the role of primary, and MySQL Router routes requests to the new primary.
- MySQL Cluster provides a high-availability, high-redundancy version of MySQL adapted for the distributed computing environment. See MySQL NDB Cluster 7.3 and NDB Cluster 7.4, which provides information about MySQL Cluster NDB 7.3 (based on MySQL 5.6 but containing the latest improvements and fixes for the `NDBCLUSTER` storage engine).

Note

The `NDBCLUSTER` storage engine is currently not supported in MySQL 8.0.

7.1 Backup and Recovery Types

This section describes the characteristics of different types of backups.

Physical (Raw) Versus Logical Backups

Physical backups consist of raw copies of the directories and files that store database contents. This type of backup is suitable for large, important databases that need to be recovered quickly when problems occur.

Logical backups save information represented as logical database structure (`CREATE DATABASE`, `CREATE TABLE` statements) and content (`INSERT` statements or delimited-text files). This type of backup is suitable for smaller amounts of data where you might edit the data values or table structure, or recreate the data on a different machine architecture.

Physical backup methods have these characteristics:

- The backup consists of exact copies of database directories and files. Typically this is a copy of all or part of the MySQL data directory.
- Physical backup methods are faster than logical because they involve only file copying without conversion.
- Output is more compact than for logical backup.
- Because backup speed and compactness are important for busy, important databases, the MySQL Enterprise Backup product performs physical backups. For an overview of the MySQL Enterprise Backup product, see Section 29.2, “MySQL Enterprise Backup Overview”.
- Backup and restore granularity ranges from the level of the entire data directory down to the level of individual files. This may or may not provide for table-level granularity, depending on storage engine. For example, `InnoDB` tables can each be in a separate file, or share file storage with other `InnoDB` tables; each `MyISAM` table corresponds uniquely to a set of files.
- In addition to databases, the backup can include any related files such as log or configuration files.

- Data from `MEMORY` tables is tricky to back up this way because their contents are not stored on disk. (The MySQL Enterprise Backup product has a feature where you can retrieve data from `MEMORY` tables during a backup.)
- Backups are portable only to other machines that have identical or similar hardware characteristics.
- Backups can be performed while the MySQL server is not running. If the server is running, it is necessary to perform appropriate locking so that the server does not change database contents during the backup. MySQL Enterprise Backup does this locking automatically for tables that require it.
- Physical backup tools include the `mysqlbackup` of MySQL Enterprise Backup for `InnoDB` or any other tables, or file system-level commands (such as `cp`, `scp`, `tar`, `rsync`) for `MyISAM` tables.
- For restore:
  - MySQL Enterprise Backup restores `InnoDB` and other tables that it backed up.
  - `ndb_restore` restores `NDB` tables.
  - Files copied at the file system level can be copied back to their original locations with file system commands.

Logical backup methods have these characteristics:

- The backup is done by querying the MySQL server to obtain database structure and content information.
- Backup is slower than physical methods because the server must access database information and convert it to logical format. If the output is written on the client side, the server must also send it to the backup program.
- Output is larger than for physical backup, particularly when saved in text format.
- Backup and restore granularity is available at the server level (all databases), database level (all tables in a particular database), or table level. This is true regardless of storage engine.
- The backup does not include log or configuration files, or other database-related files that are not part of databases.
- Backups stored in logical format are machine independent and highly portable.
- Logical backups are performed with the MySQL server running. The server is not taken offline.
- Logical backup tools include the `mysqldump` program and the `SELECT ... INTO OUTFILE` statement. These work for any storage engine, even `MEMORY`.
- To restore logical backups, SQL-format dump files can be processed using the `mysql` client. To load delimited-text files, use the `LOAD DATA INFILE` statement or the `mysqlimport` client.

Online Versus Offline Backups

Online backups take place while the MySQL server is running so that the database information can be obtained from the server. Offline backups take place while the server is stopped. This distinction can also be described as “hot” versus “cold” backups; a “warm” backup is one where the server remains running but locked against modifying data while you access database files externally.

Online backup methods have these characteristics:

- The backup is less intrusive to other clients, which can connect to the MySQL server during the backup and may be able to access data depending on what operations they need to perform.

- Care must be taken to impose appropriate locking so that data modifications do not take place that would compromise backup integrity. The MySQL Enterprise Backup product does such locking automatically.

Offline backup methods have these characteristics:

- Clients can be affected adversely because the server is unavailable during backup. For that reason, such backups are often taken from a replication slave server that can be taken offline without harming availability.
- The backup procedure is simpler because there is no possibility of interference from client activity.

A similar distinction between online and offline applies for recovery operations, and similar characteristics apply. However, it is more likely that clients will be affected for online recovery than for online backup because recovery requires stronger locking. During backup, clients might be able to read data while it is being backed up. Recovery modifies data and does not just read it, so clients must be prevented from accessing data while it is being restored.

Local Versus Remote Backups

A local backup is performed on the same host where the MySQL server runs, whereas a remote backup is done from a different host. For some types of backups, the backup can be initiated from a remote host even if the output is written locally on the server. host.

- `mysqldump` can connect to local or remote servers. For SQL output (`CREATE` and `INSERT` statements), local or remote dumps can be done and generate output on the client. For delimited-text output (with the `--tab` option), data files are created on the server host.
- `SELECT ... INTO OUTFILE` can be initiated from a local or remote client host, but the output file is created on the server host.
- Physical backup methods typically are initiated locally on the MySQL server host so that the server can be taken offline, although the destination for copied files might be remote.

Snapshot Backups

Some file system implementations enable “snapshots” to be taken. These provide logical copies of the file system at a given point in time, without requiring a physical copy of the entire file system. (For example, the implementation may use copy-on-write techniques so that only parts of the file system modified after the snapshot time need be copied.) MySQL itself does not provide the capability for taking file system snapshots. It is available through third-party solutions such as Veritas, LVM, or ZFS.

Full Versus Incremental Backups

A full backup includes all data managed by a MySQL server at a given point in time. An incremental backup consists of the changes made to the data during a given time span (from one point in time to another). MySQL has different ways to perform full backups, such as those described earlier in this section. Incremental backups are made possible by enabling the server's binary log, which the server uses to record data changes.

Full Versus Point-in-Time (Incremental) Recovery

A full recovery restores all data from a full backup. This restores the server instance to the state that it had when the backup was made. If that state is not sufficiently current, a full recovery can be followed by recovery of incremental backups made since the full backup, to bring the server to a more up-to-date state.

Incremental recovery is recovery of changes made during a given time span. This is also called point-in-time recovery because it makes a server's state current up to a given time. Point-in-time recovery is based on the binary log and typically follows a full recovery from the backup files that restores the

server to its state when the backup was made. Then the data changes written in the binary log files are applied as incremental recovery to redo data modifications and bring the server up to the desired point in time.

Table Maintenance

Data integrity can be compromised if tables become corrupt. For `InnoDB` tables, this is not a typical issue. For programs to check `MyISAM` tables and repair them if problems are found, see Section 7.6, “MyISAM Table Maintenance and Crash Recovery”.

Backup Scheduling, Compression, and Encryption

Backup scheduling is valuable for automating backup procedures. Compression of backup output reduces space requirements, and encryption of the output provides better security against unauthorized access of backed-up data. MySQL itself does not provide these capabilities. The MySQL Enterprise Backup product can compress `InnoDB` backups, and compression or encryption of backup output can be achieved using file system utilities. Other third-party solutions may be available.

7.2 Database Backup Methods

This section summarizes some general methods for making backups.

Making a Hot Backup with MySQL Enterprise Backup

Customers of MySQL Enterprise Edition can use the MySQL Enterprise Backup product to do physical backups of entire instances or selected databases, tables, or both. This product includes features for incremental and compressed backups. Backing up the physical database files makes restore much faster than logical techniques such as the `mysqldump` command. `InnoDB` tables are copied using a hot backup mechanism. (Ideally, the `InnoDB` tables should represent a substantial majority of the data.) Tables from other storage engines are copied using a warm backup mechanism. For an overview of the MySQL Enterprise Backup product, see Section 29.2, “MySQL Enterprise Backup Overview”.

Making Backups with mysqldump

The `mysqldump` program can make backups. It can back up all kinds of tables. (See Section 7.4, “Using mysqldump for Backups”.)

For `InnoDB` tables, it is possible to perform an online backup that takes no locks on tables using the `--single-transaction` option to `mysqldump`. See Section 7.3.1, “Establishing a Backup Policy”.

Making Delimited-Text File Backups

To create a text file containing a table's data, you can use `SELECT * INTO OUTFILE 'file_name' FROM tbl_name`. The file is created on the MySQL server host, not the client host. For this statement, the output file cannot already exist because permitting files to be overwritten constitutes a security risk. See Section 13.2.10, “SELECT Syntax”. This method works for any kind of data file, but saves only table data, not the table structure.

Another way to create text data files (along with files containing `CREATE TABLE` statements for the backed up tables) is to use `mysqldump` with the `--tab` option. See Section 7.4.3, “Dumping Data in Delimited-Text Format with mysqldump”.

To reload a delimited-text data file, use `LOAD DATA INFILE` or `mysqlimport`.

Making Incremental Backups by Enabling the Binary Log

MySQL supports incremental backups: You must start the server with the `--log-bin` option to enable binary logging; see Section 5.4.4, “The Binary Log”. The binary log files provide you with the

information you need to replicate changes to the database that are made subsequent to the point at which you performed a backup. At the moment you want to make an incremental backup (containing all changes that happened since the last full or incremental backup), you should rotate the binary log by using `FLUSH LOGS`. This done, you need to copy to the backup location all binary logs which range from the one of the moment of the last full or incremental backup to the last but one. These binary logs are the incremental backup; at restore time, you apply them as explained in Section 7.5, "Point-in-Time (Incremental) Recovery Using the Binary Log". The next time you do a full backup, you should also rotate the binary log using `FLUSH LOGS` or `mysqldump --flush-logs`. See Section 4.5.4, "`mysqldump` — A Database Backup Program".

Making Backups Using Replication Slaves

If you have performance problems with your master server while making backups, one strategy that can help is to set up replication and perform backups on the slave rather than on the master. See Section 18.3.1, "Using Replication for Backups".

If you are backing up a slave replication server, you should back up its master info and relay log info repositories (see Section 18.2.4, "Replication Relay and Status Logs") when you back up the slave's databases, regardless of the backup method you choose. These information files are always needed to resume replication after you restore the slave's data. If your slave is replicating `LOAD DATA INFILE` statements, you should also back up any `SQL_LOAD-*` files that exist in the directory that the slave uses for this purpose. The slave needs these files to resume replication of any interrupted `LOAD DATA INFILE` operations. The location of this directory is the value of the `--slave-load-tmpdir` option. If the server was not started with that option, the directory location is the value of the `tmpdir` system variable.

Recovering Corrupt Tables

If you have to restore `MyISAM` tables that have become corrupt, try to recover them using `REPAIR TABLE` or `myisamchk -r` first. That should work in 99.9% of all cases. If `myisamchk` fails, see Section 7.6, "MyISAM Table Maintenance and Crash Recovery".

Making Backups Using a File System Snapshot

If you are using a Veritas file system, you can make a backup like this:

1. From a client program, execute `FLUSH TABLES WITH READ LOCK`.
2. From another shell, execute `mount vxfs snapshot`.
3. From the first client, execute `UNLOCK TABLES`.
4. Copy files from the snapshot.
5. Unmount the snapshot.

Similar snapshot capabilities may be available in other file systems, such as LVM or ZFS.

7.3 Example Backup and Recovery Strategy

This section discusses a procedure for performing backups that enables you to recover data after several types of crashes:

- Operating system crash
- Power failure
- File system crash
- Hardware problem (hard drive, motherboard, and so forth)

The example commands do not include options such as `--user` and `--password` for the `mysqldump` and `mysql` client programs. You should include such options as necessary to enable client programs to connect to the MySQL server.

Assume that data is stored in the `InnoDB` storage engine, which has support for transactions and automatic crash recovery. Assume also that the MySQL server is under load at the time of the crash. If it were not, no recovery would ever be needed.

For cases of operating system crashes or power failures, we can assume that MySQL's disk data is available after a restart. The `InnoDB` data files might not contain consistent data due to the crash, but `InnoDB` reads its logs and finds in them the list of pending committed and noncommitted transactions that have not been flushed to the data files. `InnoDB` automatically rolls back those transactions that were not committed, and flushes to its data files those that were committed. Information about this recovery process is conveyed to the user through the MySQL error log. The following is an example log excerpt:

```
InnoDB: Database was not shut down normally.
InnoDB: Starting recovery from log files...
InnoDB: Starting log scan based on checkpoint at
InnoDB: log sequence number 0 13674004
InnoDB: Doing recovery: scanned up to log sequence number 0 13739520
InnoDB: Doing recovery: scanned up to log sequence number 0 13805056
InnoDB: Doing recovery: scanned up to log sequence number 0 13870592
InnoDB: Doing recovery: scanned up to log sequence number 0 13936128
...
InnoDB: Doing recovery: scanned up to log sequence number 0 20555264
InnoDB: Doing recovery: scanned up to log sequence number 0 20620800
InnoDB: Doing recovery: scanned up to log sequence number 0 20664692
InnoDB: 1 uncommitted transaction(s) which must be rolled back
InnoDB: Starting rollback of uncommitted transactions
InnoDB: Rolling back trx no 16745
InnoDB: Rolling back of trx no 16745 completed
InnoDB: Rollback of uncommitted transactions completed
InnoDB: Starting an apply batch of log records to the database...
InnoDB: Apply batch completed
InnoDB: Started
mysqld: ready for connections
```

For the cases of file system crashes or hardware problems, we can assume that the MySQL disk data is *not* available after a restart. This means that MySQL fails to start successfully because some blocks of disk data are no longer readable. In this case, it is necessary to reformat the disk, install a new one, or otherwise correct the underlying problem. Then it is necessary to recover our MySQL data from backups, which means that backups must already have been made. To make sure that is the case, design and implement a backup policy.

7.3.1 Establishing a Backup Policy

To be useful, backups must be scheduled regularly. A full backup (a snapshot of the data at a point in time) can be done in MySQL with several tools. For example, MySQL Enterprise Backup can perform a physical backup of an entire instance, with optimizations to minimize overhead and avoid disruption when backing up `InnoDB` data files; `mysqldump` provides online logical backup. This discussion uses `mysqldump`.

Assume that we make a full backup of all our `InnoDB` tables in all databases using the following command on Sunday at 1 p.m., when load is low:

```
shell> mysqldump --all-databases --master-data --single-transaction > backup_sunday_1_PM.sql
```

The resulting `.sql` file produced by `mysqldump` contains a set of SQL `INSERT` statements that can be used to reload the dumped tables at a later time.

This backup operation acquires a global read lock on all tables at the beginning of the dump (using `FLUSH TABLES WITH READ LOCK`). As soon as this lock has been acquired, the binary log

coordinates are read and the lock is released. If long updating statements are running when the `FLUSH` statement is issued, the backup operation may stall until those statements finish. After that, the dump becomes lock-free and does not disturb reads and writes on the tables.

It was assumed earlier that the tables to back up are `InnoDB` tables, so `--single-transaction` uses a consistent read and guarantees that data seen by `mysqldump` does not change. (Changes made by other clients to `InnoDB` tables are not seen by the `mysqldump` process.) If the backup operation includes nontransactional tables, consistency requires that they do not change during the backup. For example, for the `MyISAM` tables in the `mysql` database, there must be no administrative changes to MySQL accounts during the backup.

Full backups are necessary, but it is not always convenient to create them. They produce large backup files and take time to generate. They are not optimal in the sense that each successive full backup includes all data, even that part that has not changed since the previous full backup. It is more efficient to make an initial full backup, and then to make incremental backups. The incremental backups are smaller and take less time to produce. The tradeoff is that, at recovery time, you cannot restore your data just by reloading the full backup. You must also process the incremental backups to recover the incremental changes.

To make incremental backups, we need to save the incremental changes. In MySQL, these changes are represented in the binary log, so the MySQL server should always be started with the `--log-bin` option to enable that log. With binary logging enabled, the server writes each data change into a file while it updates data. Looking at the data directory of a MySQL server that was started with the `--log-bin` option and that has been running for some days, we find these MySQL binary log files:

```
-rw-rw---- 1 guilhem  guilhem   1277324 Nov 10 23:59 gbichot2-bin.000001
-rw-rw---- 1 guilhem  guilhem         4 Nov 10 23:59 gbichot2-bin.000002
-rw-rw---- 1 guilhem  guilhem        79 Nov 11 11:06 gbichot2-bin.000003
-rw-rw---- 1 guilhem  guilhem       508 Nov 11 11:08 gbichot2-bin.000004
-rw-rw---- 1 guilhem  guilhem 220047446 Nov 12 16:47 gbichot2-bin.000005
-rw-rw---- 1 guilhem  guilhem    998412 Nov 14 10:08 gbichot2-bin.000006
-rw-rw---- 1 guilhem  guilhem       361 Nov 14 10:07 gbichot2-bin.index
```

Each time it restarts, the MySQL server creates a new binary log file using the next number in the sequence. While the server is running, you can also tell it to close the current binary log file and begin a new one manually by issuing a `FLUSH LOGS` SQL statement or with a `mysqladmin flush-logs` command. `mysqldump` also has an option to flush the logs. The `.index` file in the data directory contains the list of all MySQL binary logs in the directory.

The MySQL binary logs are important for recovery because they form the set of incremental backups. If you make sure to flush the logs when you make your full backup, the binary log files created afterward contain all the data changes made since the backup. Let's modify the previous `mysqldump` command a bit so that it flushes the MySQL binary logs at the moment of the full backup, and so that the dump file contains the name of the new current binary log:

```
shell> mysqldump --single-transaction --flush-logs --master-data=2 \
         --all-databases > backup_sunday_1_PM.sql
```

After executing this command, the data directory contains a new binary log file, `gbichot2-bin.000007`, because the `--flush-logs` option causes the server to flush its logs. The `--master-data` option causes `mysqldump` to write binary log information to its output, so the resulting `.sql` dump file includes these lines:

```
-- Position to start replication or point-in-time recovery from
-- CHANGE MASTER TO MASTER_LOG_FILE='gbichot2-bin.000007',MASTER_LOG_POS=4;
```

Because the `mysqldump` command made a full backup, those lines mean two things:

- The dump file contains all changes made before any changes written to the `gbichot2-bin.000007` binary log file or higher.

- All data changes logged after the backup are not present in the dump file, but are present in the `gbichot2-bin.000007` binary log file or higher.

On Monday at 1 p.m., we can create an incremental backup by flushing the logs to begin a new binary log file. For example, executing a `mysqladmin flush-logs` command creates `gbichot2-bin.000008`. All changes between the Sunday 1 p.m. full backup and Monday 1 p.m. will be in the `gbichot2-bin.000007` file. This incremental backup is important, so it is a good idea to copy it to a safe place. (For example, back it up on tape or DVD, or copy it to another machine.) On Tuesday at 1 p.m., execute another `mysqladmin flush-logs` command. All changes between Monday 1 p.m. and Tuesday 1 p.m. will be in the `gbichot2-bin.000008` file (which also should be copied somewhere safe).

The MySQL binary logs take up disk space. To free up space, purge them from time to time. One way to do this is by deleting the binary logs that are no longer needed, such as when we make a full backup:

```
shell> mysqldump --single-transaction --flush-logs --master-data=2 \
         --all-databases --delete-master-logs > backup_sunday_1_PM.sql
```

Note

Deleting the MySQL binary logs with `mysqldump --delete-master-logs` can be dangerous if your server is a replication master server, because slave servers might not yet fully have processed the contents of the binary log. The description for the `PURGE BINARY LOGS` statement explains what should be verified before deleting the MySQL binary logs. See Section 13.4.1.1, “PURGE BINARY LOGS Syntax”.

7.3.2 Using Backups for Recovery

Now, suppose that we have a catastrophic crash on Wednesday at 8 a.m. that requires recovery from backups. To recover, first we restore the last full backup we have (the one from Sunday 1 p.m.). The full backup file is just a set of SQL statements, so restoring it is very easy:

```
shell> mysql < backup_sunday_1_PM.sql
```

At this point, the data is restored to its state as of Sunday 1 p.m.. To restore the changes made since then, we must use the incremental backups; that is, the `gbichot2-bin.000007` and `gbichot2-bin.000008` binary log files. Fetch the files if necessary from where they were backed up, and then process their contents like this:

```
shell> mysqlbinlog gbichot2-bin.000007 gbichot2-bin.000008 | mysql
```

We now have recovered the data to its state as of Tuesday 1 p.m., but still are missing the changes from that date to the date of the crash. To not lose them, we would have needed to have the MySQL server store its MySQL binary logs into a safe location (RAID disks, SAN, ...) different from the place where it stores its data files, so that these logs were not on the destroyed disk. (That is, we can start the server with a `--log-bin` option that specifies a location on a different physical device from the one on which the data directory resides. That way, the logs are safe even if the device containing the directory is lost.) If we had done this, we would have the `gbichot2-bin.000009` file (and any subsequent files) at hand, and we could apply them using `mysqlbinlog` and `mysql` to restore the most recent data changes with no loss up to the moment of the crash:

```
shell> mysqlbinlog gbichot2-bin.000009 ... | mysql
```

For more information about using `mysqlbinlog` to process binary log files, see Section 7.5, “Point-in-Time (Incremental) Recovery Using the Binary Log”.

7.3.3 Backup Strategy Summary

In case of an operating system crash or power failure, `InnoDB` itself does all the job of recovering data. But to make sure that you can sleep well, observe the following guidelines:

- Always run the MySQL server with the `--log-bin` option, or even `--log-bin=log_name`, where the log file name is located on some safe media different from the drive on which the data directory is located. If you have such safe media, this technique can also be good for disk load balancing (which results in a performance improvement).
- Make periodic full backups, using the `mysqldump` command shown earlier in Section 7.3.1, “Establishing a Backup Policy”, that makes an online, nonblocking backup.
- Make periodic incremental backups by flushing the logs with `FLUSH LOGS` or `mysqladmin flush-logs`.

7.4 Using mysqldump for Backups

This section describes how to use `mysqldump` to produce dump files, and how to reload dump files. A dump file can be used in several ways:

- As a backup to enable data recovery in case of data loss.
- As a source of data for setting up replication slaves.
- As a source of data for experimentation:
  - To make a copy of a database that you can use without changing the original data.
  - To test potential upgrade incompatibilities.

`mysqldump` produces two types of output, depending on whether the `--tab` option is given:

- Without `--tab`, `mysqldump` writes SQL statements to the standard output. This output consists of `CREATE` statements to create dumped objects (databases, tables, stored routines, and so forth), and `INSERT` statements to load data into tables. The output can be saved in a file and reloaded later using `mysql` to recreate the dumped objects. Options are available to modify the format of the SQL statements, and to control which objects are dumped.
- With `--tab`, `mysqldump` produces two output files for each dumped table. The server writes one file as tab-delimited text, one line per table row. This file is named *`tbl_name`*`.txt` in the output directory. The server also sends a `CREATE TABLE` statement for the table to `mysqldump`, which writes it as a file named *`tbl_name`*`.sql` in the output directory.

7.4.1 Dumping Data in SQL Format with mysqldump

This section describes how to use `mysqldump` to create SQL-format dump files. For information about reloading such dump files, see Section 7.4.2, “Reloading SQL-Format Backups”.

By default, `mysqldump` writes information as SQL statements to the standard output. You can save the output in a file:

```
shell> mysqldump [arguments] > file_name
```

To dump all databases, invoke `mysqldump` with the `--all-databases` option:

```
shell> mysqldump --all-databases > dump.sql
```

To dump only specific databases, name them on the command line and use the `--databases` option:

```
shell> mysqldump --databases db1 db2 db3 > dump.sql
```

The `--databases` option causes all names on the command line to be treated as database names. Without this option, `mysqldump` treats the first name as a database name and those following as table names.

With `--all-databases` or `--databases`, `mysqldump` writes `CREATE DATABASE` and `USE` statements prior to the dump output for each database. This ensures that when the dump file is reloaded, it creates each database if it does not exist and makes it the default database so database contents are loaded into the same database from which they came. If you want to cause the dump file to force a drop of each database before recreating it, use the `--add-drop-database` option as well. In this case, `mysqldump` writes a `DROP DATABASE` statement preceding each `CREATE DATABASE` statement.

To dump a single database, name it on the command line:

```
shell> mysqldump --databases test > dump.sql
```

In the single-database case, it is permissible to omit the `--databases` option:

```
shell> mysqldump test > dump.sql
```

The difference between the two preceding commands is that without `--databases`, the dump output contains no `CREATE DATABASE` or `USE` statements. This has several implications:

- When you reload the dump file, you must specify a default database name so that the server knows which database to reload.
- For reloading, you can specify a database name different from the original name, which enables you to reload the data into a different database.
- If the database to be reloaded does not exist, you must create it first.
- Because the output will contain no `CREATE DATABASE` statement, the `--add-drop-database` option has no effect. If you use it, it produces no `DROP DATABASE` statement.

To dump only specific tables from a database, name them on the command line following the database name:

```
shell> mysqldump test t1 t3 t7 > dump.sql
```

7.4.2 Reloading SQL-Format Backups

To reload a dump file written by `mysqldump` that consists of SQL statements, use it as input to the `mysql` client. If the dump file was created by `mysqldump` with the `--all-databases` or `--databases` option, it contains `CREATE DATABASE` and `USE` statements and it is not necessary to specify a default database into which to load the data:

```
shell> mysql < dump.sql
```

Alternatively, from within `mysql`, use a `source` command:

```
mysql> source dump.sql
```

If the file is a single-database dump not containing `CREATE DATABASE` and `USE` statements, create the database first (if necessary):

```
shell> mysqladmin create db1
```

Then specify the database name when you load the dump file:

```
shell> mysql db1 < dump.sql
```

Alternatively, from within `mysql`, create the database, select it as the default database, and load the dump file:

```
mysql> CREATE DATABASE IF NOT EXISTS db1;
mysql> USE db1;
mysql> source dump.sql
```

Note

For Windows PowerShell users: Because the "<" character is reserved for future use in PowerShell, an alternative approach is required, such as using quotes `cmd.exe /c "mysql < dump.sql"`.

7.4.3 Dumping Data in Delimited-Text Format with mysqldump

This section describes how to use `mysqldump` to create delimited-text dump files. For information about reloading such dump files, see Section 7.4.4, “Reloading Delimited-Text Format Backups”.

If you invoke `mysqldump` with the `--tab=dir_name` option, it uses *`dir_name`* as the output directory and dumps tables individually in that directory using two files for each table. The table name is the base name for these files. For a table named `t1`, the files are named `t1.sql` and `t1.txt`. The `.sql` file contains a `CREATE TABLE` statement for the table. The `.txt` file contains the table data, one line per table row.

The following command dumps the contents of the `db1` database to files in the `/tmp` database:

```
shell> mysqldump --tab=/tmp db1
```

The `.txt` files containing table data are written by the server, so they are owned by the system account used for running the server. The server uses `SELECT ... INTO OUTFILE` to write the files, so you must have the `FILE` privilege to perform this operation, and an error occurs if a given `.txt` file already exists.

The server sends the `CREATE` definitions for dumped tables to `mysqldump`, which writes them to `.sql` files. These files therefore are owned by the user who executes `mysqldump`.

It is best that `--tab` be used only for dumping a local server. If you use it with a remote server, the `--tab` directory must exist on both the local and remote hosts, and the `.txt` files will be written by the server in the remote directory (on the server host), whereas the `.sql` files will be written by `mysqldump` in the local directory (on the client host).

For `mysqldump --tab`, the server by default writes table data to `.txt` files one line per row with tabs between column values, no quotation marks around column values, and newline as the line terminator. (These are the same defaults as for `SELECT ... INTO OUTFILE`.)

To enable data files to be written using a different format, `mysqldump` supports these options:

- `--fields-terminated-by=`*`str`*

  The string for separating column values (default: tab).

- `--fields-enclosed-by=`*`char`*

  The character within which to enclose column values (default: no character).

- `--fields-optionally-enclosed-by=char`

  The character within which to enclose non-numeric column values (default: no character).

- `--fields-escaped-by=char`

  The character for escaping special characters (default: no escaping).

- `--lines-terminated-by=str`

  The line-termination string (default: newline).

Depending on the value you specify for any of these options, it might be necessary on the command line to quote or escape the value appropriately for your command interpreter. Alternatively, specify the value using hex notation. Suppose that you want `mysqldump` to quote column values within double quotation marks. To do so, specify double quote as the value for the `--fields-enclosed-by` option. But this character is often special to command interpreters and must be treated specially. For example, on Unix, you can quote the double quote like this:

```
--fields-enclosed-by='"'
```

On any platform, you can specify the value in hex:

```
--fields-enclosed-by=0x22
```

It is common to use several of the data-formatting options together. For example, to dump tables in comma-separated values format with lines terminated by carriage-return/newline pairs (`\r\n`), use this command (enter it on a single line):

```
shell> mysqldump --tab=/tmp --fields-terminated-by=,
         --fields-enclosed-by='"' --lines-terminated-by=0x0d0a db1
```

Should you use any of the data-formatting options to dump table data, you will need to specify the same format when you reload data files later, to ensure proper interpretation of the file contents.

7.4.4 Reloading Delimited-Text Format Backups

For backups produced with `mysqldump --tab`, each table is represented in the output directory by an `.sql` file containing the `CREATE TABLE` statement for the table, and a `.txt` file containing the table data. To reload a table, first change location into the output directory. Then process the `.sql` file with `mysql` to create an empty table and process the `.txt` file to load the data into the table:

```
shell> mysql db1 < t1.sql
shell> mysqlimport db1 t1.txt
```

An alternative to using `mysqlimport` to load the data file is to use the `LOAD DATA INFILE` statement from within the `mysql` client:

```
mysql> USE db1;
mysql> LOAD DATA INFILE 't1.txt' INTO TABLE t1;
```

If you used any data-formatting options with `mysqldump` when you initially dumped the table, you must use the same options with `mysqlimport` or `LOAD DATA INFILE` to ensure proper interpretation of the data file contents:

```
shell> mysqlimport --fields-terminated-by=,
         --fields-enclosed-by='"' --lines-terminated-by=0x0d0a db1 t1.txt
```

Or:

```
mysql> USE db1;
mysql> LOAD DATA INFILE 't1.txt' INTO TABLE t1
    -> FIELDS TERMINATED BY ',' FIELDS ENCLOSED BY '"'
    -> LINES TERMINATED BY '\r\n';
```

7.4.5 mysqldump Tips

This section surveys techniques that enable you to use `mysqldump` to solve specific problems:

- How to make a copy a database
- How to copy a database from one server to another
- How to dump stored programs (stored procedures and functions, triggers, and events)
- How to dump definitions and data separately

7.4.5.1 Making a Copy of a Database

```
shell> mysqldump db1 > dump.sql
shell> mysqladmin create db2
shell> mysql db2 < dump.sql
```

Do not use `--databases` on the `mysqldump` command line because that causes `USE db1` to be included in the dump file, which overrides the effect of naming `db2` on the `mysql` command line.

7.4.5.2 Copy a Database from one Server to Another

On Server 1:

```
shell> mysqldump --databases db1 > dump.sql
```

Copy the dump file from Server 1 to Server 2.

On Server 2:

```
shell> mysql < dump.sql
```

Use of `--databases` with the `mysqldump` command line causes the dump file to include `CREATE DATABASE` and `USE` statements that create the database if it does exist and make it the default database for the reloaded data.

Alternatively, you can omit `--databases` from the `mysqldump` command. Then you will need to create the database on Server 2 (if necessary) and specify it as the default database when you reload the dump file.

On Server 1:

```
shell> mysqldump db1 > dump.sql
```

On Server 2:

```
shell> mysqladmin create db1
shell> mysql db1 < dump.sql
```

You can specify a different database name in this case, so omitting `--databases` from the `mysqldump` command enables you to dump data from one database and load it into another.

7.4.5.3 Dumping Stored Programs

Several options control how `mysqldump` handles stored programs (stored procedures and functions, triggers, and events):

- `--events`: Dump Event Scheduler events
- `--routines`: Dump stored procedures and functions
- `--triggers`: Dump triggers for tables

The `--triggers` option is enabled by default so that when tables are dumped, they are accompanied by any triggers they have. The other options are disabled by default and must be specified explicitly to dump the corresponding objects. To disable any of these options explicitly, use its skip form: `--skip-events`, `--skip-routines`, or `--skip-triggers`.

7.4.5.4 Dumping Table Definitions and Content Separately

The `--no-data` option tells `mysqldump` not to dump table data, resulting in the dump file containing only statements to create the tables. Conversely, the `--no-create-info` option tells `mysqldump` to suppress `CREATE` statements from the output, so that the dump file contains only table data.

For example, to dump table definitions and data separately for the `test` database, use these commands:

```
shell> mysqldump --no-data test > dump-defs.sql
shell> mysqldump --no-create-info test > dump-data.sql
```

For a definition-only dump, add the `--routines` and `--events` options to also include stored routine and event definitions:

```
shell> mysqldump --no-data --routines --events test > dump-defs.sql
```

7.4.5.5 Using mysqldump to Test for Upgrade Incompatibilities

When contemplating a MySQL upgrade, it is prudent to install the newer version separately from your current production version. Then you can dump the database and database object definitions from the production server and load them into the new server to verify that they are handled properly. (This is also useful for testing downgrades.)

On the production server:

```
shell> mysqldump --all-databases --no-data --routines --events > dump-defs.sql
```

On the upgraded server:

```
shell> mysql < dump-defs.sql
```

Because the dump file does not contain table data, it can be processed quickly. This enables you to spot potential incompatibilities without waiting for lengthy data-loading operations. Look for warnings or errors while the dump file is being processed.

After you have verified that the definitions are handled properly, dump the data and try to load it into the upgraded server.

On the production server:

```
shell> mysqldump --all-databases --no-create-info > dump-data.sql
```

On the upgraded server:

```
shell> mysql < dump-data.sql
```

Now check the table contents and run some test queries.

7.5 Point-in-Time (Incremental) Recovery Using the Binary Log

Point-in-time recovery refers to recovery of data changes made since a given point in time. Typically, this type of recovery is performed after restoring a full backup that brings the server to its state as of the time the backup was made. (The full backup can be made in several ways, such as those listed in Section 7.2, “Database Backup Methods”.) Point-in-time recovery then brings the server up to date incrementally from the time of the full backup to a more recent time.

Note

Many of the examples here use the `mysql` client to process binary log output produced by `mysqlbinlog`. If your binary log contains `\0` (null) characters, that output cannot be parsed by `mysql` unless you invoke it with the `--binary-mode` option.

Point-in-time recovery is based on these principles:

- The source of information for point-in-time recovery is the set of incremental backups represented by the binary log files generated subsequent to the full backup operation. Therefore, the server must be started with the `--log-bin` option to enable binary logging (see Section 5.4.4, “The Binary Log”).

  To restore data from the binary log, you must know the name and location of the current binary log files. By default, the server creates binary log files in the data directory, but a path name can be specified with the `--log-bin` option to place the files in a different location. Section 5.4.4, “The Binary Log”.

  To see a listing of all binary log files, use this statement:

  ```
  mysql> SHOW BINARY LOGS;
  ```

  To determine the name of the current binary log file, issue the following statement:

  ```
  mysql> SHOW MASTER STATUS;
  ```

- The `mysqlbinlog` utility converts the events in the binary log files from binary format to text so that they can be executed or viewed. `mysqlbinlog` has options for selecting sections of the binary log based on event times or position of events within the log. See Section 4.6.8, “`mysqlbinlog` — Utility for Processing Binary Log Files”.

- Executing events from the binary log causes the data modifications they represent to be redone. This enables recovery of data changes for a given span of time. To execute events from the binary log, process `mysqlbinlog` output using the `mysql` client:

  ```
  shell> mysqlbinlog binlog_files | mysql -u root -p
  ```

- Viewing log contents can be useful when you need to determine event times or positions to select partial log contents prior to executing events. To view events from the log, send `mysqlbinlog` output into a paging program:

  ```
  shell> mysqlbinlog binlog_files | more
  ```

  Alternatively, save the output in a file and view the file in a text editor:

```
shell> mysqlbinlog binlog_files > tmpfile
shell> ... edit tmpfile ...
```

- Saving the output in a file is useful as a preliminary to executing the log contents with certain events removed, such as an accidental `DROP DATABASE`. You can delete from the file any statements not to be executed before executing its contents. After editing the file, execute the contents as follows:

```
shell> mysql -u root -p < tmpfile
```

If you have more than one binary log to execute on the MySQL server, the safe method is to process them all using a single connection to the server. Here is an example that demonstrates what may be *unsafe*:

```
shell> mysqlbinlog binlog.000001 | mysql -u root -p # DANGER!!
shell> mysqlbinlog binlog.000002 | mysql -u root -p # DANGER!!
```

Processing binary logs this way using different connections to the server causes problems if the first log file contains a `CREATE TEMPORARY TABLE` statement and the second log contains a statement that uses the temporary table. When the first `mysql` process terminates, the server drops the temporary table. When the second `mysql` process attempts to use the table, the server reports "unknown table."

To avoid problems like this, use a *single* connection to execute the contents of all binary logs that you want to process. Here is one way to do so:

```
shell> mysqlbinlog binlog.000001 binlog.000002 | mysql -u root -p
```

Another approach is to write all the logs to a single file and then process the file:

```
shell> mysqlbinlog binlog.000001 >  /tmp/statements.sql
shell> mysqlbinlog binlog.000002 >> /tmp/statements.sql
shell> mysql -u root -p -e "source /tmp/statements.sql"
```

When writing to a dump file while reading back from a binary log containing GTIDs (see Section 18.1.3, "Replication with Global Transaction Identifiers"), use the `--skip-gtids` option with `mysqlbinlog`, like this:

```
shell> mysqlbinlog --skip-gtids binlog.000001 >  /tmp/dump.sql
shell> mysqlbinlog --skip-gtids binlog.000002 >> /tmp/dump.sql
shell> mysql -u root -p -e "source /tmp/dump.sql"
```

7.5.1 Point-in-Time Recovery Using Event Times

To indicate the start and end times for recovery, specify the `--start-datetime` and `--stop-datetime` options for `mysqlbinlog`, in `DATETIME` format. As an example, suppose that exactly at 10:00 a.m. on April 20, 2005 an SQL statement was executed that deleted a large table. To restore the table and data, you could restore the previous night's backup, and then execute the following command:

```
shell> mysqlbinlog --stop-datetime="2005-04-20 9:59:59" \
        /var/log/mysql/bin.123456 | mysql -u root -p
```

This command recovers all of the data up until the date and time given by the `--stop-datetime` option. If you did not detect the erroneous SQL statement that was entered until hours later, you will probably also want to recover the activity that occurred afterward. Based on this, you could run `mysqlbinlog` again with a start date and time, like so:

```
shell> mysqlbinlog --start-datetime="2005-04-20 10:01:00" \
        /var/log/mysql/bin.123456 | mysql -u root -p
```

In this command, the SQL statements logged from 10:01 a.m. on will be re-executed. The combination of restoring of the previous night's dump file and the two `mysqlbinlog` commands restores everything up until one second before 10:00 a.m. and everything from 10:01 a.m. on.

To use this method of point-in-time recovery, you should examine the log to be sure of the exact times to specify for the commands. To display the log file contents without executing them, use this command:

```
shell> mysqlbinlog /var/log/mysql/bin.123456 > /tmp/mysql_restore.sql
```

Then open the `/tmp/mysql_restore.sql` file with a text editor to examine it.

Excluding specific changes by specifying times for `mysqlbinlog` does not work well if multiple statements executed at the same time as the one to be excluded.

7.5.2 Point-in-Time Recovery Using Event Positions

Instead of specifying dates and times, the `--start-position` and `--stop-position` options for `mysqlbinlog` can be used for specifying log positions. They work the same as the start and stop date options, except that you specify log position numbers rather than dates. Using positions may enable you to be more precise about which part of the log to recover, especially if many transactions occurred around the same time as a damaging SQL statement. To determine the position numbers, run `mysqlbinlog` for a range of times near the time when the unwanted transaction was executed, but redirect the results to a text file for examination. This can be done like so:

```
shell> mysqlbinlog --start-datetime="2005-04-20 9:55:00" \
        --stop-datetime="2005-04-20 10:05:00" \
        /var/log/mysql/bin.123456 > /tmp/mysql_restore.sql
```

This command creates a small text file in the `/tmp` directory that contains the SQL statements around the time that the deleterious SQL statement was executed. Open this file with a text editor and look for the statement that you do not want to repeat. Determine the positions in the binary log for stopping and resuming the recovery and make note of them. Positions are labeled as `log_pos` followed by a number. After restoring the previous backup file, use the position numbers to process the binary log file. For example, you would use commands something like these:

```
shell> mysqlbinlog --stop-position=368312 /var/log/mysql/bin.123456 \
        | mysql -u root -p

shell> mysqlbinlog --start-position=368315 /var/log/mysql/bin.123456 \
        | mysql -u root -p
```

The first command recovers all the transactions up until the stop position given. The second command recovers all transactions from the starting position given until the end of the binary log. Because the output of `mysqlbinlog` includes `SET TIMESTAMP` statements before each SQL statement recorded, the recovered data and related MySQL logs will reflect the original times at which the transactions were executed.

7.6 MyISAM Table Maintenance and Crash Recovery

This section discusses how to use `myisamchk` to check or repair `MyISAM` tables (tables that have `.MYD` and `.MYI` files for storing data and indexes). For general `myisamchk` background, see Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”. Other table-repair information can be found at Section 2.10.3, “Rebuilding or Repairing Tables or Indexes”.

You can use `myisamchk` to check, repair, or optimize database tables. The following sections describe how to perform these operations and how to set up a table maintenance schedule. For information

about using myisamchk to get information about your tables, see Section 4.6.4.5, “Obtaining Table Information with myisamchk”.

Even though table repair with myisamchk is quite secure, it is always a good idea to make a backup *before* doing a repair or any maintenance operation that could make a lot of changes to a table.

myisamchk operations that affect indexes can cause MyISAM FULLTEXT indexes to be rebuilt with full-text parameters that are incompatible with the values used by the MySQL server. To avoid this problem, follow the guidelines in Section 4.6.4.1, “myisamchk General Options”.

MyISAM table maintenance can also be done using the SQL statements that perform operations similar to what myisamchk can do:

- To check MyISAM tables, use CHECK TABLE.
- To repair MyISAM tables, use REPAIR TABLE.
- To optimize MyISAM tables, use OPTIMIZE TABLE.
- To analyze MyISAM tables, use ANALYZE TABLE.

For additional information about these statements, see Section 13.7.2, “Table Maintenance Statements”.

These statements can be used directly or by means of the mysqlcheck client program. One advantage of these statements over myisamchk is that the server does all the work. With myisamchk, you must make sure that the server does not use the tables at the same time so that there is no unwanted interaction between myisamchk and the server.

7.6.1 Using myisamchk for Crash Recovery

This section describes how to check for and deal with data corruption in MySQL databases. If your tables become corrupted frequently, you should try to find the reason why. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.

For an explanation of how MyISAM tables can become corrupted, see Section 16.2.4, “MyISAM Table Problems”.

If you run mysqld with external locking disabled (which is the default), you cannot reliably use myisamchk to check a table when mysqld is using the same table. If you can be certain that no one will access the tables through mysqld while you run myisamchk, you only have to execute mysqladmin flush-tables before you start checking the tables. If you cannot guarantee this, you must stop mysqld while you check the tables. If you run myisamchk to check tables that mysqld is updating at the same time, you may get a warning that a table is corrupt even when it is not.

If the server is run with external locking enabled, you can use myisamchk to check tables at any time. In this case, if the server tries to update a table that myisamchk is using, the server will wait for myisamchk to finish before it continues.

If you use myisamchk to repair or optimize tables, you *must* always ensure that the mysqld server is not using the table (this also applies if external locking is disabled). If you do not stop mysqld, you should at least do a mysqladmin flush-tables before you run myisamchk. Your tables *may become corrupted* if the server and myisamchk access the tables simultaneously.

When performing crash recovery, it is important to understand that each MyISAM table *tbl_name* in a database corresponds to the three files in the database directory shown in the following table.

File	Purpose
tbl_name.MYD	Data file
tbl_name.MYI	Index file

Each of these three file types is subject to corruption in various ways, but problems occur most often in data files and index files.

`myisamchk` works by creating a copy of the `.MYD` data file row by row. It ends the repair stage by removing the old `.MYD` file and renaming the new file to the original file name. If you use `--quick`, `myisamchk` does not create a temporary `.MYD` file, but instead assumes that the `.MYD` file is correct and generates only a new index file without touching the `.MYD` file. This is safe, because `myisamchk` automatically detects whether the `.MYD` file is corrupt and aborts the repair if it is. You can also specify the `--quick` option twice to `myisamchk`. In this case, `myisamchk` does not abort on some errors (such as duplicate-key errors) but instead tries to resolve them by modifying the `.MYD` file. Normally the use of two `--quick` options is useful only if you have too little free disk space to perform a normal repair. In this case, you should at least make a backup of the table before running `myisamchk`.

7.6.2 How to Check MyISAM Tables for Errors

To check a `MyISAM` table, use the following commands:

- `myisamchk *tbl_name*`

  This finds 99.99% of all errors. What it cannot find is corruption that involves *only* the data file (which is very unusual). If you want to check a table, you should normally run `myisamchk` without options or with the `-s` (silent) option.

- `myisamchk -m *tbl_name*`

  This finds 99.999% of all errors. It first checks all index entries for errors and then reads through all rows. It calculates a checksum for all key values in the rows and verifies that the checksum matches the checksum for the keys in the index tree.

- `myisamchk -e *tbl_name*`

  This does a complete and thorough check of all data (`-e` means “extended check”). It does a check-read of every key for each row to verify that they indeed point to the correct row. This may take a long time for a large table that has many indexes. Normally, `myisamchk` stops after the first error it finds. If you want to obtain more information, you can add the `-v` (verbose) option. This causes `myisamchk` to keep going, up through a maximum of 20 errors.

- `myisamchk -e -i *tbl_name*`

  This is like the previous command, but the `-i` option tells `myisamchk` to print additional statistical information.

In most cases, a simple `myisamchk` command with no arguments other than the table name is sufficient to check a table.

7.6.3 How to Repair MyISAM Tables

The discussion in this section describes how to use `myisamchk` on `MyISAM` tables (extensions `.MYI` and `.MYD`).

You can also use the `CHECK TABLE` and `REPAIR TABLE` statements to check and repair `MyISAM` tables. See Section 13.7.2.2, “CHECK TABLE Syntax”, and Section 13.7.2.5, “REPAIR TABLE Syntax”.

Symptoms of corrupted tables include queries that abort unexpectedly and observable errors such as these:

- Can't find file *`tbl_name`*`.MYI` (Errcode: *`nnn`*)
- Unexpected end of file

- Record file is crashed
- Got error *nnn* from table handler

To get more information about the error, run `perror nnn`, where *nnn* is the error number. The following example shows how to use `perror` to find the meanings for the most common error numbers that indicate a problem with a table:

```
shell> perror 126 127 132 134 135 136 141 144 145
MySQL error code 126 = Index file is crashed
MySQL error code 127 = Record-file is crashed
MySQL error code 132 = Old database file
MySQL error code 134 = Record was already deleted (or record file crashed)
MySQL error code 135 = No more room in record file
MySQL error code 136 = No more room in index file
MySQL error code 141 = Duplicate unique key or constraint on write or update
MySQL error code 144 = Table is crashed and last repair failed
MySQL error code 145 = Table was marked as crashed and should be repaired
```

Note that error 135 (no more room in record file) and error 136 (no more room in index file) are not errors that can be fixed by a simple repair. In this case, you must use `ALTER TABLE` to increase the `MAX_ROWS` and `AVG_ROW_LENGTH` table option values:

```
ALTER TABLE tbl_name MAX_ROWS=xxx AVG_ROW_LENGTH=yyy;
```

If you do not know the current table option values, use `SHOW CREATE TABLE`.

For the other errors, you must repair your tables. `myisamchk` can usually detect and fix most problems that occur.

The repair process involves up to three stages, described here. Before you begin, you should change location to the database directory and check the permissions of the table files. On Unix, make sure that they are readable by the user that `mysqld` runs as (and to you, because you need to access the files you are checking). If it turns out you need to modify files, they must also be writable by you.

This section is for the cases where a table check fails (such as those described in Section 7.6.2, “How to Check MyISAM Tables for Errors”), or you want to use the extended features that `myisamchk` provides.

The `myisamchk` options used for table maintenance with are described in Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”. `myisamchk` also has variables that you can set to control memory allocation that may improve performance. See Section 4.6.4.6, “myisamchk Memory Usage”.

If you are going to repair a table from the command line, you must first stop the `mysqld` server. Note that when you do `mysqladmin shutdown` on a remote server, the `mysqld` server is still available for a while after `mysqladmin` returns, until all statement-processing has stopped and all index changes have been flushed to disk.

Stage 1: Checking your tables

Run `myisamchk *.MYI` or `myisamchk -e *.MYI` if you have more time. Use the `-s` (silent) option to suppress unnecessary information.

If the `mysqld` server is stopped, you should use the `--update-state` option to tell `myisamchk` to mark the table as “checked.”

You have to repair only those tables for which `myisamchk` announces an error. For such tables, proceed to Stage 2.

If you get unexpected errors when checking (such as `out of memory` errors), or if `myisamchk` crashes, go to Stage 3.

Stage 2: Easy safe repair

First, try `myisamchk -r -q tbl_name` (`-r -q` means “quick recovery mode”). This attempts to repair the index file without touching the data file. If the data file contains everything that it should and the delete links point at the correct locations within the data file, this should work, and the table is fixed. Start repairing the next table. Otherwise, use the following procedure:

1. Make a backup of the data file before continuing.
2. Use `myisamchk -r tbl_name` (`-r` means “recovery mode”). This removes incorrect rows and deleted rows from the data file and reconstructs the index file.
3. If the preceding step fails, use `myisamchk --safe-recover tbl_name`. Safe recovery mode uses an old recovery method that handles a few cases that regular recovery mode does not (but is slower).

Note

If you want a repair operation to go much faster, you should set the values of the `sort_buffer_size` and `key_buffer_size` variables each to about 25% of your available memory when running `myisamchk`.

If you get unexpected errors when repairing (such as `out of memory` errors), or if `myisamchk` crashes, go to Stage 3.

Stage 3: Difficult repair

You should reach this stage only if the first 16KB block in the index file is destroyed or contains incorrect information, or if the index file is missing. In this case, it is necessary to create a new index file. Do so as follows:

1. Move the data file to a safe place.
2. Use the table description file to create new (empty) data and index files:

```
shell> mysql db_name
mysql> SET autocommit=1;
mysql> TRUNCATE TABLE tbl_name;
mysql> quit
```

3. Copy the old data file back onto the newly created data file. (Do not just move the old file back onto the new file. You want to retain a copy in case something goes wrong.)

Important

If you are using replication, you should stop it prior to performing the above procedure, since it involves file system operations, and these are not logged by MySQL.

Go back to Stage 2. `myisamchk -r -q` should work. (This should not be an endless loop.)

You can also use the `REPAIR TABLE tbl_name USE_FRM` SQL statement, which performs the whole procedure automatically. There is also no possibility of unwanted interaction between a utility and the server, because the server does all the work when you use `REPAIR TABLE`. See Section 13.7.2.5, “REPAIR TABLE Syntax”.

7.6.4 MyISAM Table Optimization

To coalesce fragmented rows and eliminate wasted space that results from deleting or updating rows, run `myisamchk` in recovery mode:

```
shell> myisamchk -r tbl_name
```

You can optimize a table in the same way by using the `OPTIMIZE TABLE` SQL statement. `OPTIMIZE TABLE` does a table repair and a key analysis, and also sorts the index tree so that key lookups are faster. There is also no possibility of unwanted interaction between a utility and the server, because the server does all the work when you use `OPTIMIZE TABLE`. See Section 13.7.2.4, “OPTIMIZE TABLE Syntax”.

`myisamchk` has a number of other options that you can use to improve the performance of a table:

- `--analyze` or `-a`: Perform key distribution analysis. This improves join performance by enabling the join optimizer to better choose the order in which to join the tables and which indexes it should use.
- `--sort-index` or `-S`: Sort the index blocks. This optimizes seeks and makes table scans that use indexes faster.
- `--sort-records=index_num` or `-R index_num`: Sort data rows according to a given index. This makes your data much more localized and may speed up range-based `SELECT` and `ORDER BY` operations that use this index.

For a full description of all available options, see Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”.

7.6.5 Setting Up a MyISAM Table Maintenance Schedule

It is a good idea to perform table checks on a regular basis rather than waiting for problems to occur. One way to check and repair `MyISAM` tables is with the `CHECK TABLE` and `REPAIR TABLE` statements. See Section 13.7.2, “Table Maintenance Statements”.

Another way to check tables is to use `myisamchk`. For maintenance purposes, you can use `myisamchk -s`. The `-s` option (short for `--silent`) causes `myisamchk` to run in silent mode, printing messages only when errors occur.

It is also a good idea to enable automatic `MyISAM` table checking. For example, whenever the machine has done a restart in the middle of an update, you usually need to check each table that could have been affected before it is used further. (These are “expected crashed tables.”) To cause the server to check `MyISAM` tables automatically, start it with the `--myisam-recover-options` option. See Section 5.1.4, “Server Command Options”.

You should also check your tables regularly during normal system operation. For example, you can run a `cron` job to check important tables once a week, using a line like this in a `crontab` file:

```
35 0 * * 0 /path/to/myisamchk --fast --silent /path/to/datadir/*/*.MYI
```

This prints out information about crashed tables so that you can examine and repair them as necessary.

To start with, execute `myisamchk -s` each night on all tables that have been updated during the last 24 hours. As you see that problems occur infrequently, you can back off the checking frequency to once a week or so.

Normally, MySQL tables need little maintenance. If you are performing many updates to `MyISAM` tables with dynamic-sized rows (tables with `VARCHAR`, `BLOB`, or `TEXT` columns) or have tables with many deleted rows you may want to defragment/reclaim space from the tables from time to time. You can do this by using `OPTIMIZE TABLE` on the tables in question. Alternatively, if you can stop the `mysqld` server for a while, change location into the data directory and use this command while the server is stopped:

```
shell> myisamchk -r -s --sort-index --myisam_sort_buffer_size=16M */*.MYI
```

Chapter 8 Optimization

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This chapter explains how to optimize MySQL performance and provides examples. Optimization involves configuring, tuning, and measuring performance, at several levels. Depending on your job role (developer, DBA, or a combination of both), you might optimize at the level of individual SQL statements, entire applications, a single database server, or multiple networked database servers. Sometimes you can be proactive and plan in advance for performance, while other times you might troubleshoot a configuration or code issue after a problem occurs. Optimizing CPU and memory usage can also improve scalability, allowing the database to handle more load without slowing down.

8.1 Optimization Overview

Database performance depends on several factors at the database level, such as tables, queries, and configuration settings. These software constructs result in CPU and I/O operations at the hardware level, which you must minimize and make as efficient as possible. As you work on database performance, you start by learning the high-level rules and guidelines for the software side, and measuring performance using wall-clock time. As you become an expert, you learn more about what happens internally, and start measuring things such as CPU cycles and I/O operations.

Typical users aim to get the best database performance out of their existing software and hardware configurations. Advanced users look for opportunities to improve the MySQL software itself, or develop their own storage engines and hardware appliances to expand the MySQL ecosystem.

- Optimizing at the Database Level

- Optimizing at the Hardware Level
- Balancing Portability and Performance

Optimizing at the Database Level

The most important factor in making a database application fast is its basic design:

- Are the tables structured properly? In particular, do the columns have the right data types, and does each table have the appropriate columns for the type of work? For example, applications that perform frequent updates often have many tables with few columns, while applications that analyze large amounts of data often have few tables with many columns.
- Are the right indexes in place to make queries efficient?
- Are you using the appropriate storage engine for each table, and taking advantage of the strengths and features of each storage engine you use? In particular, the choice of a transactional storage engine such as `InnoDB` or a nontransactional one such as `MyISAM` can be very important for performance and scalability.

Note

`InnoDB` is the default storage engine for new tables. In practice, the advanced `InnoDB` performance features mean that `InnoDB` tables often outperform the simpler `MyISAM` tables, especially for a busy database.

- Does each table use an appropriate row format? This choice also depends on the storage engine used for the table. In particular, compressed tables use less disk space and so require less disk I/O to read and write the data. Compression is available for all kinds of workloads with `InnoDB` tables, and for read-only `MyISAM` tables.
- Does the application use an appropriate locking strategy? For example, by allowing shared access when possible so that database operations can run concurrently, and requesting exclusive access when appropriate so that critical operations get top priority. Again, the choice of storage engine is significant. The `InnoDB` storage engine handles most locking issues without involvement from you, allowing for better concurrency in the database and reducing the amount of experimentation and tuning for your code.
- Are all memory areas used for caching sized correctly? That is, large enough to hold frequently accessed data, but not so large that they overload physical memory and cause paging. The main memory areas to configure are the `InnoDB` buffer pool, the `MyISAM` key cache, and the MySQL query cache.

Optimizing at the Hardware Level

Any database application eventually hits hardware limits as the database becomes more and more busy. A DBA must evaluate whether it is possible to tune the application or reconfigure the server to avoid these bottlenecks, or whether more hardware resources are required. System bottlenecks typically arise from these sources:

- Disk seeks. It takes time for the disk to find a piece of data. With modern disks, the mean time for this is usually lower than 10ms, so we can in theory do about 100 seeks a second. This time improves slowly with new disks and is very hard to optimize for a single table. The way to optimize seek time is to distribute the data onto more than one disk.
- Disk reading and writing. When the disk is at the correct position, we need to read or write the data. With modern disks, one disk delivers at least 10–20MB/s throughput. This is easier to optimize than seeks because you can read in parallel from multiple disks.
- CPU cycles. When the data is in main memory, we must process it to get our result. Having large tables compared to the amount of memory is the most common limiting factor. But with small tables, speed is usually not the problem.

- Memory bandwidth. When the CPU needs more data than can fit in the CPU cache, main memory bandwidth becomes a bottleneck. This is an uncommon bottleneck for most systems, but one to be aware of.

Balancing Portability and Performance

To use performance-oriented SQL extensions in a portable MySQL program, you can wrap MySQL-specific keywords in a statement within `/*! */` comment delimiters. Other SQL servers ignore the commented keywords. For information about writing comments, see Section 9.6, “Comment Syntax”.

8.2 Optimizing SQL Statements

The core logic of a database application is performed through SQL statements, whether issued directly through an interpreter or submitted behind the scenes through an API. The tuning guidelines in this section help to speed up all kinds of MySQL applications. The guidelines cover SQL operations that read and write data, the behind-the-scenes overhead for SQL operations in general, and operations used in specific scenarios such as database monitoring.

8.2.1 Optimizing SELECT Statements

Queries, in the form of `SELECT` statements, perform all the lookup operations in the database. Tuning these statements is a top priority, whether to achieve sub-second response times for dynamic web pages, or to chop hours off the time to generate huge overnight reports.

Besides `SELECT` statements, the tuning techniques for queries also apply to constructs such as `CREATE TABLE...AS SELECT`, `INSERT INTO...SELECT`, and `WHERE` clauses in `DELETE` statements. Those statements have additional performance considerations because they combine write operations with the read-oriented query operations.

The main considerations for optimizing queries are:

- To make a slow `SELECT ... WHERE` query faster, the first thing to check is whether you can add an index. Set up indexes on columns used in the `WHERE` clause, to speed up evaluation, filtering, and the final retrieval of results. To avoid wasted disk space, construct a small set of indexes that speed up many related queries used in your application.

  Indexes are especially important for queries that reference different tables, using features such as joins and foreign keys. You can use the `EXPLAIN` statement to determine which indexes are used for a `SELECT`. See Section 8.3.1, “How MySQL Uses Indexes” and Section 8.8.1, “Optimizing Queries with EXPLAIN”.

- Isolate and tune any part of the query, such as a function call, that takes excessive time. Depending on how the query is structured, a function could be called once for every row in the result set, or even once for every row in the table, greatly magnifying any inefficiency.

- Minimize the number of full table scans in your queries, particularly for big tables.

- Keep table statistics up to date by using the `ANALYZE TABLE` statement periodically, so the optimizer has the information needed to construct an efficient execution plan.

- Learn the tuning techniques, indexing techniques, and configuration parameters that are specific to the storage engine for each table. Both `InnoDB` and `MyISAM` have sets of guidelines for enabling and sustaining high performance in queries. For details, see Section 8.5.6, “Optimizing InnoDB Queries” and Section 8.6.1, “Optimizing MyISAM Queries”.

- You can optimize single-query transactions for `InnoDB` tables, using the technique in Section 8.5.3, “Optimizing InnoDB Read-Only Transactions”.

- Avoid transforming the query in ways that make it hard to understand, especially if the optimizer does some of the same transformations automatically.

- If a performance issue is not easily solved by one of the basic guidelines, investigate the internal details of the specific query by reading the `EXPLAIN` plan and adjusting your indexes, `WHERE` clauses, join clauses, and so on. (When you reach a certain level of expertise, reading the `EXPLAIN` plan might be your first step for every query.)
- Adjust the size and properties of the memory areas that MySQL uses for caching. With efficient use of the `InnoDB` buffer pool, `MyISAM` key cache, and the MySQL query cache, repeated queries run faster because the results are retrieved from memory the second and subsequent times.
- Even for a query that runs fast using the cache memory areas, you might still optimize further so that they require less cache memory, making your application more scalable. Scalability means that your application can handle more simultaneous users, larger requests, and so on without experiencing a big drop in performance.
- Deal with locking issues, where the speed of your query might be affected by other sessions accessing the tables at the same time.

8.2.1.1 WHERE Clause Optimization

This section discusses optimizations that can be made for processing `WHERE` clauses. The examples use `SELECT` statements, but the same optimizations apply for `WHERE` clauses in `DELETE` and `UPDATE` statements.

Note

Because work on the MySQL optimizer is ongoing, not all of the optimizations that MySQL performs are documented here.

You might be tempted to rewrite your queries to make arithmetic operations faster, while sacrificing readability. Because MySQL does similar optimizations automatically, you can often avoid this work, and leave the query in a more understandable and maintainable form. Some of the optimizations performed by MySQL follow:

- Removal of unnecessary parentheses:

```
   ((a AND b) AND c OR (((a AND b) AND (c AND d))))
-> (a AND b AND c) OR (a AND b AND c AND d)
```

- Constant folding:

```
   (a<b AND b=c) AND a=5
-> b>5 AND b=c AND a=5
```

- Constant condition removal (needed because of constant folding):

```
   (B>=5 AND B=5) OR (B=6 AND 5=5) OR (B=7 AND 5=6)
-> B=5 OR B=6
```

- Constant expressions used by indexes are evaluated only once.
- `COUNT(*)` on a single table without a `WHERE` is retrieved directly from the table information for `MyISAM` and `MEMORY` tables. This is also done for any `NOT NULL` expression when used with only one table.
- Early detection of invalid constant expressions. MySQL quickly detects that some `SELECT` statements are impossible and returns no rows.
- `HAVING` is merged with `WHERE` if you do not use `GROUP BY` or aggregate functions (`COUNT()`, `MIN()`, and so on).
- For each table in a join, a simpler `WHERE` is constructed to get a fast `WHERE` evaluation for the table and also to skip rows as soon as possible.

- All constant tables are read first before any other tables in the query. A constant table is any of the following:

  - An empty table or a table with one row.
  - A table that is used with a `WHERE` clause on a `PRIMARY KEY` or a `UNIQUE` index, where all index parts are compared to constant expressions and are defined as `NOT NULL`.

  All of the following tables are used as constant tables:

```
SELECT * FROM t WHERE primary_key=1;
SELECT * FROM t1,t2
  WHERE t1.primary_key=1 AND t2.primary_key=t1.id;
```

- The best join combination for joining the tables is found by trying all possibilities. If all columns in `ORDER BY` and `GROUP BY` clauses come from the same table, that table is preferred first when joining.
- If there is an `ORDER BY` clause and a different `GROUP BY` clause, or if the `ORDER BY` or `GROUP BY` contains columns from tables other than the first table in the join queue, a temporary table is created.
- If you use the `SQL_SMALL_RESULT` modifier, MySQL uses an in-memory temporary table.
- Each table index is queried, and the best index is used unless the optimizer believes that it is more efficient to use a table scan. At one time, a scan was used based on whether the best index spanned more than 30% of the table, but a fixed percentage no longer determines the choice between using an index or a scan. The optimizer now is more complex and bases its estimate on additional factors such as table size, number of rows, and I/O block size.
- In some cases, MySQL can read rows from the index without even consulting the data file. If all columns used from the index are numeric, only the index tree is used to resolve the query.
- Before each row is output, those that do not match the `HAVING` clause are skipped.

Some examples of queries that are very fast:

```
SELECT COUNT(*) FROM tbl_name;

SELECT MIN(key_part1),MAX(key_part1) FROM tbl_name;

SELECT MAX(key_part2) FROM tbl_name
  WHERE key_part1=constant;

SELECT ... FROM tbl_name
  ORDER BY key_part1,key_part2,... LIMIT 10;

SELECT ... FROM tbl_name
  ORDER BY key_part1 DESC, key_part2 DESC, ... LIMIT 10;
```

MySQL resolves the following queries using only the index tree, assuming that the indexed columns are numeric:

```
SELECT key_part1,key_part2 FROM tbl_name WHERE key_part1=val;

SELECT COUNT(*) FROM tbl_name
  WHERE key_part1=val1 AND key_part2=val2;

SELECT key_part2 FROM tbl_name GROUP BY key_part1;
```

The following queries use indexing to retrieve the rows in sorted order without a separate sorting pass:

```
SELECT ... FROM tbl_name
  ORDER BY key_part1,key_part2,... ;
```

```
SELECT ... FROM tbl_name
  ORDER BY key_part1 DESC, key_part2 DESC, ... ;
```

8.2.1.2 Range Optimization

The `range` access method uses a single index to retrieve a subset of table rows that are contained within one or several index value intervals. It can be used for a single-part or multiple-part index. The following sections give descriptions of conditions under which the optimizer uses range access.

- The Range Access Method for Single-Part Indexes
- The Range Access Method for Multiple-Part Indexes
- Equality Range Optimization of Many-Valued Comparisons
- Limiting Memory Use for Range Optimization
- Range Optimization of Row Constructor Expressions

The Range Access Method for Single-Part Indexes

For a single-part index, index value intervals can be conveniently represented by corresponding conditions in the `WHERE` clause, denoted as *range conditions* rather than "intervals."

The definition of a range condition for a single-part index is as follows:

- For both `BTREE` and `HASH` indexes, comparison of a key part with a constant value is a range condition when using the `=`, `<=>`, `IN()`, `IS NULL`, or `IS NOT NULL` operators.
- Additionally, for `BTREE` indexes, comparison of a key part with a constant value is a range condition when using the `>`, `<`, `>=`, `<=`, `BETWEEN`, `!=`, or `<>` operators, or `LIKE` comparisons if the argument to `LIKE` is a constant string that does not start with a wildcard character.
- For all index types, multiple range conditions combined with `OR` or `AND` form a range condition.

"Constant value" in the preceding descriptions means one of the following:

- A constant from the query string
- A column of a `const` or `system` table from the same join
- The result of an uncorrelated subquery
- Any expression composed entirely from subexpressions of the preceding types

Here are some examples of queries with range conditions in the `WHERE` clause:

```
SELECT * FROM t1
  WHERE key_col > 1
  AND key_col < 10;

SELECT * FROM t1
  WHERE key_col = 1
  OR key_col IN (15,18,20);

SELECT * FROM t1
  WHERE key_col LIKE 'ab%'
  OR key_col BETWEEN 'bar' AND 'foo';
```

Some nonconstant values may be converted to constants during the optimizer constant propagation phase.

MySQL tries to extract range conditions from the `WHERE` clause for each of the possible indexes. During the extraction process, conditions that cannot be used for constructing the range condition are

dropped, conditions that produce overlapping ranges are combined, and conditions that produce empty ranges are removed.

Consider the following statement, where `key1` is an indexed column and `nonkey` is not indexed:

```
SELECT * FROM t1 WHERE
  (key1 < 'abc' AND (key1 LIKE 'abcde%' OR key1 LIKE '%b')) OR
  (key1 < 'bar' AND nonkey = 4) OR
  (key1 < 'uux' AND key1 > 'z');
```

The extraction process for key `key1` is as follows:

1. Start with original `WHERE` clause:

   ```
   (key1 < 'abc' AND (key1 LIKE 'abcde%' OR key1 LIKE '%b')) OR
   (key1 < 'bar' AND nonkey = 4) OR
   (key1 < 'uux' AND key1 > 'z')
   ```

2. Remove `nonkey = 4` and `key1 LIKE '%b'` because they cannot be used for a range scan. The correct way to remove them is to replace them with `TRUE`, so that we do not miss any matching rows when doing the range scan. Having replaced them with `TRUE`, we get:

   ```
   (key1 < 'abc' AND (key1 LIKE 'abcde%' OR TRUE)) OR
   (key1 < 'bar' AND TRUE) OR
   (key1 < 'uux' AND key1 > 'z')
   ```

3. Collapse conditions that are always true or false:

   - `(key1 LIKE 'abcde%' OR TRUE)` is always true
   - `(key1 < 'uux' AND key1 > 'z')` is always false

   Replacing these conditions with constants, we get:

   ```
   (key1 < 'abc' AND TRUE) OR (key1 < 'bar' AND TRUE) OR (FALSE)
   ```

   Removing unnecessary `TRUE` and `FALSE` constants, we obtain:

   ```
   (key1 < 'abc') OR (key1 < 'bar')
   ```

4. Combining overlapping intervals into one yields the final condition to be used for the range scan:

   ```
   (key1 < 'bar')
   ```

In general (and as demonstrated by the preceding example), the condition used for a range scan is less restrictive than the `WHERE` clause. MySQL performs an additional check to filter out rows that satisfy the range condition but not the full `WHERE` clause.

The range condition extraction algorithm can handle nested `AND`/`OR` constructs of arbitrary depth, and its output does not depend on the order in which conditions appear in `WHERE` clause.

MySQL does not support merging multiple ranges for the `range` access method for spatial indexes. To work around this limitation, you can use a `UNION` with identical `SELECT` statements, except that you put each spatial predicate in a different `SELECT`.

The Range Access Method for Multiple-Part Indexes

Range conditions on a multiple-part index are an extension of range conditions for a single-part index. A range condition on a multiple-part index restricts index rows to lie within one or several key tuple intervals. Key tuple intervals are defined over a set of key tuples, using ordering from the index.

For example, consider a multiple-part index defined as `key1(key_part1, key_part2, key_part3)`, and the following set of key tuples listed in key order:

```
key_part1  key_part2  key_part3
  NULL       1          'abc'
  NULL       1          'xyz'
  NULL       2          'foo'
   1         1          'abc'
   1         1          'xyz'
   1         2          'abc'
   2         1          'aaa'
```

The condition `key_part1 = 1` defines this interval:

```
(1,-inf,-inf) <= (key_part1,key_part2,key_part3) < (1,+inf,+inf)
```

The interval covers the 4th, 5th, and 6th tuples in the preceding data set and can be used by the range access method.

By contrast, the condition `key_part3 = 'abc'` does not define a single interval and cannot be used by the range access method.

The following descriptions indicate how range conditions work for multiple-part indexes in greater detail.

- For `HASH` indexes, each interval containing identical values can be used. This means that the interval can be produced only for conditions in the following form:

  ```
     key_part1 cmp const1
  AND key_part2 cmp const2
  AND ...
  AND key_partN cmp constN;
  ```

  Here, *`const1`*, *`const2`*, ... are constants, *`cmp`* is one of the `=`, `<=>`, or `IS NULL` comparison operators, and the conditions cover all index parts. (That is, there are *`N`* conditions, one for each part of an *`N`*-part index.) For example, the following is a range condition for a three-part `HASH` index:

  ```
  key_part1 = 1 AND key_part2 IS NULL AND key_part3 = 'foo'
  ```

  For the definition of what is considered to be a constant, see The Range Access Method for Single-Part Indexes.

- For a `BTREE` index, an interval might be usable for conditions combined with `AND`, where each condition compares a key part with a constant value using `=`, `<=>`, `IS NULL`, `>`, `<`, `>=`, `<=`, `!=`, `<>`, `BETWEEN`, or `LIKE 'pattern'` (where `'pattern'` does not start with a wildcard). An interval can be used as long as it is possible to determine a single key tuple containing all rows that match the condition (or two intervals if `<>` or `!=` is used).

  The optimizer attempts to use additional key parts to determine the interval as long as the comparison operator is `=`, `<=>`, or `IS NULL`. If the operator is `>`, `<`, `>=`, `<=`, `!=`, `<>`, `BETWEEN`, or `LIKE`, the optimizer uses it but considers no more key parts. For the following expression, the optimizer uses `=` from the first comparison. It also uses `>=` from the second comparison but considers no further key parts and does not use the third comparison for interval construction:

  ```
  key_part1 = 'foo' AND key_part2 >= 10 AND key_part3 > 10
  ```

  The single interval is:

  ```
  ('foo',10,-inf) < (key_part1,key_part2,key_part3) < ('foo',+inf,+inf)
  ```

It is possible that the created interval contains more rows than the initial condition. For example, the preceding interval includes the value `('foo', 11, 0)`, which does not satisfy the original condition.

- If conditions that cover sets of rows contained within intervals are combined with `OR`, they form a condition that covers a set of rows contained within the union of their intervals. If the conditions are combined with `AND`, they form a condition that covers a set of rows contained within the intersection of their intervals. For example, for this condition on a two-part index:

```
(key_part1 = 1 AND key_part2 < 2) OR (key_part1 > 5)
```

The intervals are:

```
(1,-inf) < (key_part1,key_part2) < (1,2)
(5,-inf) < (key_part1,key_part2)
```

In this example, the interval on the first line uses one key part for the left bound and two key parts for the right bound. The interval on the second line uses only one key part. The `key_len` column in the `EXPLAIN` output indicates the maximum length of the key prefix used.

In some cases, `key_len` may indicate that a key part was used, but that might be not what you would expect. Suppose that *`key_part1`* and *`key_part2`* can be `NULL`. Then the `key_len` column displays two key part lengths for the following condition:

```
key_part1 >= 1 AND key_part2 < 2
```

But, in fact, the condition is converted to this:

```
key_part1 >= 1 AND key_part2 IS NOT NULL
```

The Range Access Method for Single-Part Indexes, describes how optimizations are performed to combine or eliminate intervals for range conditions on a single-part index. Analogous steps are performed for range conditions on multiple-part indexes.

Equality Range Optimization of Many-Valued Comparisons

Consider these expressions, where *`col_name`* is an indexed column:

```
col_name IN(val1, ..., valN)
col_name = val1 OR ... OR col_name = valN
```

Each expression is true if *`col_name`* is equal to any of several values. These comparisons are equality range comparisons (where the “range” is a single value). The optimizer estimates the cost of reading qualifying rows for equality range comparisons as follows:

- If there is a unique index on *`col_name`*, the row estimate for each range is 1 because at most one row can have the given value.
- Otherwise, any index on *`col_name`* is nonunique and the optimizer can estimate the row count for each range using dives into the index or index statistics.

With index dives, the optimizer makes a dive at each end of a range and uses the number of rows in the range as the estimate. For example, the expression *`col_name`* `IN (10, 20, 30)` has three equality ranges and the optimizer makes two dives per range to generate a row estimate. Each pair of dives yields an estimate of the number of rows that have the given value.

Index dives provide accurate row estimates, but as the number of comparison values in the expression increases, the optimizer takes longer to generate a row estimate. Use of index statistics is less accurate than index dives but permits faster row estimation for large value lists.

The `eq_range_index_dive_limit` system variable enables you to configure the number of values at which the optimizer switches from one row estimation strategy to the other. To permit use of index dives for comparisons of up to *`N`* equality ranges, set `eq_range_index_dive_limit` to *`N`* + 1. To disable use of statistics and always use index dives regardless of *`N`*, set `eq_range_index_dive_limit` to 0.

To update table index statistics for best estimates, use `ANALYZE TABLE`.

Limiting Memory Use for Range Optimization

To control the memory available to the range optimizer, use the `range_optimizer_max_mem_size` system variable:

- A value of 0 means “no limit.”
- With a value greater than 0, the optimizer tracks the memory consumed when considering the range access method. If the specified limit is about to be exceeded, the range access method is abandoned and other methods, including a full table scan, are considered instead. This could be less optimal. If this happens, the following warning occurs (where *`N`* is the current `range_optimizer_max_mem_size` value):

```
Warning    3170    Memory capacity of N bytes for
                   'range_optimizer_max_mem_size' exceeded. Range
                   optimization was not done for this query.
```

For individual queries that exceed the available range optimization memory and for which the optimizer falls back to less optimal plans, increasing the `range_optimizer_max_mem_size` value may improve performance.

To estimate the amount of memory needed to process a range expression, use these guidelines:

- For a simple query such as the following, where there is one candidate key for the range access method, each predicate combined with `OR` uses approximately 230 bytes:

```
SELECT COUNT(*) FROM t
WHERE a=1 OR a=2 OR a=3 OR .. . a=N;
```

- Similarly for a query such as the following, each predicate combined with `AND` uses approximately 125 bytes:

```
SELECT COUNT(*) FROM t
WHERE a=1 AND b=1 AND c=1 ... N;
```

- For a query with `IN()` predicates:

```
SELECT COUNT(*) FROM t
WHERE a IN (1,2, ..., M) AND b IN (1,2, ..., N);
```

  Each literal value in an `IN()` list counts as a predicate combined with `OR`. If there are two `IN()` lists, the number of predicates combined with `OR` is the product of the number of literal values in each list. Thus, the number of predicates combined with `OR` in the preceding case is *`M`* × *`N`*.

Range Optimization of Row Constructor Expressions

The optimizer is able to apply the range scan access method to queries of this form:

```
SELECT ... FROM t1 WHERE ( col_1, col_2 ) IN (( 'a', 'b' ), ( 'c', 'd' ));
```

Previously, for range scans to be used, it was necessary to write the query as:

```
SELECT ... FROM t1 WHERE ( col_1 = 'a' AND col_2 = 'b' )
OR ( col_1 = 'c' AND col_2 = 'd' );
```

For the optimizer to use a range scan, queries must satisfy these conditions:

- Only `IN()` predicates are used, not `NOT IN()`.
- On the left side of the `IN()` predicate, the row constructor contains only column references.
- On the right side of the `IN()` predicate, row constructors contain only runtime constants, which are either literals or local column references that are bound to constants during execution.
- On the right side of the `IN()` predicate, there is more than one row constructor.

For more information about the optimizer and row constructors, see Section 8.2.1.18, “Row Constructor Expression Optimization”

8.2.1.3 Index Merge Optimization

The *Index Merge* access method retrieves rows with multiple `range` scans and merges their results into one. This access method merges index scans from a single table; it does not merge scans across multiple tables. The merge can produce unions, intersections, or unions-of-intersections of its underlying scans.

Example queries for which Index Merge may be used:

```
SELECT * FROM tbl_name WHERE key1 = 10 OR key2 = 20;

SELECT * FROM tbl_name
  WHERE (key1 = 10 OR key2 = 20) AND non_key = 30;

SELECT * FROM t1, t2
  WHERE (t1.key1 IN (1,2) OR t1.key2 LIKE 'value%')
  AND t2.key1 = t1.some_col;

SELECT * FROM t1, t2
  WHERE t1.key1 = 1
  AND (t2.key1 = t1.some_col OR t2.key2 = t1.some_col2);
```

In `EXPLAIN` output, the Index Merge method appears as `index_merge` in the `type` column. In this case, the `key` column contains a list of indexes used, and `key_len` contains a list of the longest key parts for those indexes.

Note

The Index Merge optimization algorithm has the following known deficiencies:

- If your query has a complex `WHERE` clause with deep `AND`/`OR` nesting and MySQL does not choose the optimal plan, try distributing terms using the following identity transformations:

  ```
  (x AND y) OR z => (x OR z) AND (y OR z)
  (x OR y) AND z => (x AND z) OR (y AND z)
  ```

- Index Merge is not applicable to full-text indexes.

The Index Merge access method has several algorithms, which are displayed in the `Extra` field of `EXPLAIN` output:

- `Using intersect(...)`
- `Using union(...)`

- `Using sort_union(...)`

The following sections describe these algorithms in greater detail. The optimizer chooses between different possible Index Merge algorithms and other access methods based on cost estimates of the various available options.

Use of Index Merge is subject to the value of the `index_merge`, `index_merge_intersection`, `index_merge_union`, and `index_merge_sort_union` flags of the `optimizer_switch` system variable. See Section 8.9.3, “Switchable Optimizations”. By default, all those flags are `on`. To enable only certain algorithms, set `index_merge` to `off`, and enable only such of the others as should be permitted. Optimizer hints may also be applied; see Optimizer Hints for Index Merge Optimization.

- The Index Merge Intersection Access Algorithm
- The Index Merge Union Access Algorithm
- The Index Merge Sort-Union Access Algorithm
- Optimizer Hints for Index Merge Optimization

The Index Merge Intersection Access Algorithm

This access algorithm is applicable when a `WHERE` clause is converted to several range conditions on different keys combined with `AND`, and each condition is one of the following:

- An *`N`*-part expression of this form, where the index has exactly *`N`* parts (that is, all index parts are covered):

```
key_part1 = const1 AND key_part2 = const2 ... AND key_partN = constN
```

- Any range condition over a primary key of an `InnoDB` table.

Examples:

```
SELECT * FROM innodb_table
  WHERE primary_key < 10 AND key_col1 = 20;

SELECT * FROM tbl_name
  WHERE (key1_part1 = 1 AND key1_part2 = 2) AND key2 = 2;
```

The Index Merge intersection algorithm performs simultaneous scans on all used indexes and produces the intersection of row sequences that it receives from the merged index scans.

If all columns used in the query are covered by the used indexes, full table rows are not retrieved (`EXPLAIN` output contains `Using index` in `Extra` field in this case). Here is an example of such a query:

```
SELECT COUNT(*) FROM t1 WHERE key1 = 1 AND key2 = 1;
```

If the used indexes do not cover all columns used in the query, full rows are retrieved only when the range conditions for all used keys are satisfied.

If one of the merged conditions is a condition over a primary key of an `InnoDB` table, it is not used for row retrieval, but is used to filter out rows retrieved using other conditions.

The Index Merge Union Access Algorithm

The criteria for this algorithm are similar to those for the Index Merge intersection algorithm. The algorithm is applicable when the table's `WHERE` clause is converted to several range conditions on different keys combined with `OR`, and each condition is one of the following:

- An *N*-part expression of this form, where the index has exactly *N* parts (that is, all index parts are covered):

```
key_part1 = const1 AND key_part2 = const2 ... AND key_partN = constN
```

- Any range condition over a primary key of an InnoDB table.
- A condition for which the Index Merge intersection algorithm is applicable.

Examples:

```
SELECT * FROM t1
  WHERE key1 = 1 OR key2 = 2 OR key3 = 3;

SELECT * FROM innodb_table
  WHERE (key1 = 1 AND key2 = 2)
     OR (key3 = 'foo' AND key4 = 'bar') AND key5 = 5;
```

The Index Merge Sort-Union Access Algorithm

This access algorithm is applicable when the WHERE clause is converted to several range conditions combined by OR, but the Index Merge union algorithm is not applicable.

Examples:

```
SELECT * FROM tbl_name
  WHERE key_col1 < 10 OR key_col2 < 20;

SELECT * FROM tbl_name
  WHERE (key_col1 > 10 OR key_col2 = 20) AND nonkey_col = 30;
```

The difference between the sort-union algorithm and the union algorithm is that the sort-union algorithm must first fetch row IDs for all rows and sort them before returning any rows.

Optimizer Hints for Index Merge Optimization

In addition to using the optimizer_switch system variable to control optimizer use of the Index Merge algorithms session-wide, MySQL supports optimizer hints to influence the optimizer on a per-statement basis. See Section 8.9.2, “Optimizer Hints”.

8.2.1.4 Engine Condition Pushdown Optimization

This optimization improves the efficiency of direct comparisons between a nonindexed column and a constant. In such cases, the condition is “pushed down” to the storage engine for evaluation. This optimization can be used only by the NDB storage engine.

Note

The NDB storage engine is currently not available in MySQL 8.0. If you are interested in using MySQL Cluster, see MySQL NDB Cluster 7.5 and NDB Cluster 7.6, which provides information about MySQL Cluster NDB 7.5 (based on MySQL 5.7 but containing the latest improvements and fixes for the NDBCLUSTER storage engine).

For MySQL Cluster, this optimization can eliminate the need to send nonmatching rows over the network between the cluster's data nodes and the MySQL server that issued the query, and can speed up queries where it is used by a factor of 5 to 10 times over cases where condition pushdown could be but is not used.

Suppose that a MySQL Cluster table is defined as follows:

```
CREATE TABLE t1 (
    a INT,
    b INT,
    KEY(a)
) ENGINE=NDB;
```

Condition pushdown can be used with queries such as the one shown here, which includes a comparison between a nonindexed column and a constant:

```
SELECT a, b FROM t1 WHERE b = 10;
```

The use of condition pushdown can be seen in the output of `EXPLAIN`:

```
mysql> EXPLAIN SELECT a,b FROM t1 WHERE b = 10\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
         type: ALL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 10
        Extra: Using where with pushed condition
```

However, condition pushdown *cannot* be used with either of these two queries:

```
SELECT a,b FROM t1 WHERE a = 10;
SELECT a,b FROM t1 WHERE b + 1 = 10;
```

Condition pushdown is not applicable to the first query because an index exists on column `a`. (An index access method would be more efficient and so would be chosen in preference to condition pushdown.) Condition pushdown cannot be employed for the second query because the comparison involving the nonindexed column `b` is indirect. (However, condition pushdown could be applied if you were to reduce `b + 1 = 10` to `b = 9` in the `WHERE` clause.)

Condition pushdown may also be employed when an indexed column is compared with a constant using a `>` or `<` operator:

```
mysql> EXPLAIN SELECT a, b FROM t1 WHERE a < 2\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
         type: range
possible_keys: a
          key: a
      key_len: 5
          ref: NULL
         rows: 2
        Extra: Using where with pushed condition
```

Other supported comparisons for condition pushdown include the following:

- *`column`* `[NOT] LIKE` *`pattern`*

  `pattern` must be a string literal containing the pattern to be matched; for syntax, see Section 12.5.1, “String Comparison Functions”.

- *`column`* `IS [NOT] NULL`

- *`column`* `IN (`*`value_list`*`)`

Each item in the *`value_list`* must be a constant, literal value.

- *`column`* `BETWEEN` *`constant1`* `AND` *`constant2`*

  `constant1` and *`constant2`* must each be a constant, literal value.

In all of the cases in the preceding list, it is possible for the condition to be converted into the form of one or more direct comparisons between a column and a constant.

Engine condition pushdown is enabled by default. To disable it at server startup, set the `optimizer_switch` system variable. For example, in a `my.cnf` file, use these lines:

```
[mysqld]
optimizer_switch=engine_condition_pushdown=off
```

At runtime, disable condition pushdown like this:

```
SET optimizer_switch='engine_condition_pushdown=off';
```

Limitations. Engine condition pushdown is subject to the following limitations:

- Condition pushdown is supported only by the `NDB` storage engine.
- Columns may be compared with constants only; however, this includes expressions which evaluate to constant values.
- Columns used in comparisons cannot be of any of the `BLOB` or `TEXT` types.
- A string value to be compared with a column must use the same collation as the column.
- Joins are not directly supported; conditions involving multiple tables are pushed separately where possible. Use extended `EXPLAIN` output to determine which conditions are actually pushed down. See Section 8.8.3, “Extended EXPLAIN Output Format”.

8.2.1.5 Index Condition Pushdown Optimization

Index Condition Pushdown (ICP) is an optimization for the case where MySQL retrieves rows from a table using an index. Without ICP, the storage engine traverses the index to locate rows in the base table and returns them to the MySQL server which evaluates the `WHERE` condition for the rows. With ICP enabled, and if parts of the `WHERE` condition can be evaluated by using only columns from the index, the MySQL server pushes this part of the `WHERE` condition down to the storage engine. The storage engine then evaluates the pushed index condition by using the index entry and only if this is satisfied is the row read from the table. ICP can reduce the number of times the storage engine must access the base table and the number of times the MySQL server must access the storage engine.

Applicability of the Index Condition Pushdown optimization is subject to these conditions:

- ICP is used for the `range`, `ref`, `eq_ref`, and `ref_or_null` access methods when there is a need to access full table rows.
- ICP can be used for `InnoDB` and `MyISAM` tables, including partitioned `InnoDB` and `MyISAM` tables.
- For `InnoDB` tables, ICP is used only for secondary indexes. The goal of ICP is to reduce the number of full-row reads and thereby reduce I/O operations. For `InnoDB` clustered indexes, the complete record is already read into the `InnoDB` buffer. Using ICP in this case does not reduce I/O.
- ICP is not supported with secondary indexes created on virtual generated columns. `InnoDB` supports secondary indexes on virtual generated columns.
- Conditions that refer to subqueries cannot be pushed down.

- Conditions that refer to stored functions cannot be pushed down. Storage engines cannot invoke stored functions.
- Triggered conditions cannot be pushed down. (For information about triggered conditions, see Section 8.2.2.4, “Optimizing Subqueries with the EXISTS Strategy”.)

To understand how this optimization works, first consider how an index scan proceeds when Index Condition Pushdown is not used:

1. Get the next row, first by reading the index tuple, and then by using the index tuple to locate and read the full table row.
2. Test the part of the `WHERE` condition that applies to this table. Accept or reject the row based on the test result.

Using Index Condition Pushdown, the scan proceeds like this instead:

1. Get the next row's index tuple (but not the full table row).
2. Test the part of the `WHERE` condition that applies to this table and can be checked using only index columns. If the condition is not satisfied, proceed to the index tuple for the next row.
3. If the condition is satisfied, use the index tuple to locate and read the full table row.
4. Test the remaining part of the `WHERE` condition that applies to this table. Accept or reject the row based on the test result.

`EXPLAIN` output shows `Using index condition` in the `Extra` column when Index Condition Pushdown is used. It does not show `Using index` because that does not apply when full table rows must be read.

Suppose that a table contains information about people and their addresses and that the table has an index defined as `INDEX (zipcode, lastname, firstname)`. If we know a person's `zipcode` value but are not sure about the last name, we can search like this:

```
SELECT * FROM people
  WHERE zipcode='95054'
  AND lastname LIKE '%etrunia%'
  AND address LIKE '%Main Street%';
```

MySQL can use the index to scan through people with `zipcode='95054'`. The second part (`lastname LIKE '%etrunia%'`) cannot be used to limit the number of rows that must be scanned, so without Index Condition Pushdown, this query must retrieve full table rows for all people who have `zipcode='95054'`.

With Index Condition Pushdown, MySQL checks the `lastname LIKE '%etrunia%'` part before reading the full table row. This avoids reading full rows corresponding to index tuples that match the `zipcode` condition but not the `lastname` condition.

Index Condition Pushdown is enabled by default. It can be controlled with the `optimizer_switch` system variable by setting the `index_condition_pushdown` flag:

```
SET optimizer_switch = 'index_condition_pushdown=off';
SET optimizer_switch = 'index_condition_pushdown=on';
```

See Section 8.9.3, “Switchable Optimizations”.

8.2.1.6 Nested-Loop Join Algorithms

MySQL executes joins between tables using a nested-loop algorithm or variations on it.

- Nested-Loop Join Algorithm

- Block Nested-Loop Join Algorithm

Nested-Loop Join Algorithm

A simple nested-loop join (NLJ) algorithm reads rows from the first table in a loop one at a time, passing each row to a nested loop that processes the next table in the join. This process is repeated as many times as there remain tables to be joined.

Assume that a join between three tables `t1`, `t2`, and `t3` is to be executed using the following join types:

```
Table   Join Type
t1      range
t2      ref
t3      ALL
```

If a simple NLJ algorithm is used, the join is processed like this:

```
for each row in t1 matching range {
  for each row in t2 matching reference key {
    for each row in t3 {
      if row satisfies join conditions, send to client
    }
  }
}
```

Because the NLJ algorithm passes rows one at a time from outer loops to inner loops, it typically reads tables processed in the inner loops many times.

Block Nested-Loop Join Algorithm

A Block Nested-Loop (BNL) join algorithm uses buffering of rows read in outer loops to reduce the number of times that tables in inner loops must be read. For example, if 10 rows are read into a buffer and the buffer is passed to the next inner loop, each row read in the inner loop can be compared against all 10 rows in the buffer. This reduces by an order of magnitude the number of times the inner table must be read.

MySQL join buffering has these characteristics:

- Join buffering can be used when the join is of type `ALL` or `index` (in other words, when no possible keys can be used, and a full scan is done, of either the data or index rows, respectively), or `range`. Use of buffering is also applicable to outer joins, as described in Section 8.2.1.11, “Block Nested-Loop and Batched Key Access Joins”.
- A join buffer is never allocated for the first nonconstant table, even if it would be of type `ALL` or `index`.
- Only columns of interest to a join are stored in its join buffer, not whole rows.
- The `join_buffer_size` system variable determines the size of each join buffer used to process a query.
- One buffer is allocated for each join that can be buffered, so a given query might be processed using multiple join buffers.
- A join buffer is allocated prior to executing the join and freed after the query is done.

For the example join described previously for the NLJ algorithm (without buffering), the join is done as follows using join buffering:

```
for each row in t1 matching range {
  for each row in t2 matching reference key {
```

```
    store used columns from t1, t2 in join buffer
    if buffer is full {
      for each row in t3 {
        for each t1, t2 combination in join buffer {
          if row satisfies join conditions, send to client
        }
      }
      empty join buffer
    }
  }
}

if buffer is not empty {
  for each row in t3 {
    for each t1, t2 combination in join buffer {
      if row satisfies join conditions, send to client
    }
  }
}
```

If *`S`* is the size of each stored `t1`, `t2` combination is the join buffer and *`C`* is the number of combinations in the buffer, the number of times table `t3` is scanned is:

```
(S * C)/join_buffer_size + 1
```

The number of `t3` scans decreases as the value of `join_buffer_size` increases, up to the point when `join_buffer_size` is large enough to hold all previous row combinations. At that point, no speed is gained by making it larger.

8.2.1.7 Nested Join Optimization

The syntax for expressing joins permits nested joins. The following discussion refers to the join syntax described in Section 13.2.10.2, “JOIN Syntax”.

The syntax of *`table_factor`* is extended in comparison with the SQL Standard. The latter accepts only *`table_reference`*, not a list of them inside a pair of parentheses. This is a conservative extension if we consider each comma in a list of *`table_reference`* items as equivalent to an inner join. For example:

```
SELECT * FROM t1 LEFT JOIN (t2, t3, t4)
                 ON (t2.a=t1.a AND t3.b=t1.b AND t4.c=t1.c)
```

Is equivalent to:

```
SELECT * FROM t1 LEFT JOIN (t2 CROSS JOIN t3 CROSS JOIN t4)
                 ON (t2.a=t1.a AND t3.b=t1.b AND t4.c=t1.c)
```

In MySQL, `CROSS JOIN` is syntactically equivalent to `INNER JOIN`; they can replace each other. In standard SQL, they are not equivalent. `INNER JOIN` is used with an `ON` clause; `CROSS JOIN` is used otherwise.

In general, parentheses can be ignored in join expressions containing only inner join operations. Consider this join expression:

```
t1 LEFT JOIN (t2 LEFT JOIN t3 ON t2.b=t3.b OR t2.b IS NULL)
   ON t1.a=t2.a
```

After removing parentheses and grouping operations to the left, that join expression transforms into this expression:

```
(t1 LEFT JOIN t2 ON t1.a=t2.a) LEFT JOIN t3
    ON t2.b=t3.b OR t2.b IS NULL
```

Yet, the two expressions are not equivalent. To see this, suppose that the tables `t1`, `t2`, and `t3` have the following state:

- Table `t1` contains rows `(1)`, `(2)`
- Table `t2` contains row `(1,101)`
- Table `t3` contains row `(101)`

In this case, the first expression returns a result set including the rows `(1,1,101,101)`, `(2,NULL,NULL,NULL)`, whereas the second expression returns the rows `(1,1,101,101)`, `(2,NULL,NULL,101)`:

```
mysql> SELECT *
       FROM t1
            LEFT JOIN
            (t2 LEFT JOIN t3 ON t2.b=t3.b OR t2.b IS NULL)
            ON t1.a=t2.a;
+------+------+------+------+
| a    | a    | b    | b    |
+------+------+------+------+
|    1 |    1 |  101 |  101 |
|    2 | NULL | NULL | NULL |
+------+------+------+------+

mysql> SELECT *
       FROM (t1 LEFT JOIN t2 ON t1.a=t2.a)
            LEFT JOIN t3
            ON t2.b=t3.b OR t2.b IS NULL;
+------+------+------+------+
| a    | a    | b    | b    |
+------+------+------+------+
|    1 |    1 |  101 |  101 |
|    2 | NULL | NULL |  101 |
+------+------+------+------+
```

In the following example, an outer join operation is used together with an inner join operation:

```
t1 LEFT JOIN (t2, t3) ON t1.a=t2.a
```

That expression cannot be transformed into the following expression:

```
t1 LEFT JOIN t2 ON t1.a=t2.a, t3
```

For the given table states, the two expressions return different sets of rows:

```
mysql> SELECT *
       FROM t1 LEFT JOIN (t2, t3) ON t1.a=t2.a;
+------+------+------+------+
| a    | a    | b    | b    |
+------+------+------+------+
|    1 |    1 |  101 |  101 |
|    2 | NULL | NULL | NULL |
+------+------+------+------+

mysql> SELECT *
       FROM t1 LEFT JOIN t2 ON t1.a=t2.a, t3;
+------+------+------+------+
| a    | a    | b    | b    |
+------+------+------+------+
|    1 |    1 |  101 |  101 |
|    2 | NULL | NULL |  101 |
+------+------+------+------+
```

Therefore, if we omit parentheses in a join expression with outer join operators, we might change the result set for the original expression.

More exactly, we cannot ignore parentheses in the right operand of the left outer join operation and in the left operand of a right join operation. In other words, we cannot ignore parentheses for the inner table expressions of outer join operations. Parentheses for the other operand (operand for the outer table) can be ignored.

The following expression:

```
(t1,t2) LEFT JOIN t3 ON P(t2.b,t3.b)
```

Is equivalent to this expression for any tables `t1,t2,t3` and any condition `P` over attributes `t2.b` and `t3.b`:

```
t1, t2 LEFT JOIN t3 ON P(t2.b,t3.b)
```

Whenever the order of execution of join operations in a join expression (*`join_table`*) is not from left to right, we talk about nested joins. Consider the following queries:

```
SELECT * FROM t1 LEFT JOIN (t2 LEFT JOIN t3 ON t2.b=t3.b) ON t1.a=t2.a
  WHERE t1.a > 1

SELECT * FROM t1 LEFT JOIN (t2, t3) ON t1.a=t2.a
  WHERE (t2.b=t3.b OR t2.b IS NULL) AND t1.a > 1
```

Those queries are considered to contain these nested joins:

```
t2 LEFT JOIN t3 ON t2.b=t3.b
t2, t3
```

In the first query, the nested join is formed with a left join operation. In the second query, it is formed with an inner join operation.

In the first query, the parentheses can be omitted: The grammatical structure of the join expression will dictate the same order of execution for join operations. For the second query, the parentheses cannot be omitted, although the join expression here can be interpreted unambiguously without them. In our extended syntax, the parentheses in `(t2, t3)` of the second query are required, although theoretically the query could be parsed without them: We still would have unambiguous syntactical structure for the query because `LEFT JOIN` and `ON` play the role of the left and right delimiters for the expression `(t2,t3)`.

The preceding examples demonstrate these points:

- For join expressions involving only inner joins (and not outer joins), parentheses can be removed and joins evaluated left to right. In fact, tables can be evaluated in any order.
- The same is not true, in general, for outer joins or for outer joins mixed with inner joins. Removal of parentheses may change the result.

Queries with nested outer joins are executed in the same pipeline manner as queries with inner joins. More exactly, a variation of the nested-loop join algorithm is exploited. Recall the algorithm by which the nested-loop join executes a query (see Section 8.2.1.6, “Nested-Loop Join Algorithms”). Suppose that a join query over 3 tables `T1,T2,T3` has this form:

```
SELECT * FROM T1 INNER JOIN T2 ON P1(T1,T2)
                 INNER JOIN T3 ON P2(T2,T3)
  WHERE P(T1,T2,T3)
```

Here, `P1(T1,T2)` and `P2(T3,T3)` are some join conditions (on expressions), whereas `P(T1,T2,T3)` is a condition over columns of tables `T1,T2,T3`.

The nested-loop join algorithm would execute this query in the following manner:

```
FOR each row t1 in T1 {
  FOR each row t2 in T2 such that P1(t1,t2) {
    FOR each row t3 in T3 such that P2(t2,t3) {
      IF P(t1,t2,t3) {
         t:=t1||t2||t3; OUTPUT t;
      }
    }
  }
}
```

The notation `t1||t2||t3` indicates a row constructed by concatenating the columns of rows `t1`, `t2`, and `t3`. In some of the following examples, `NULL` where a table name appears means a row in which `NULL` is used for each column of that table. For example, `t1||t2||NULL` indicates a row constructed by concatenating the columns of rows `t1` and `t2`, and `NULL` for each column of `t3`. Such a row is said to be `NULL`-complemented.

Now consider a query with nested outer joins:

```
SELECT * FROM T1 LEFT JOIN
             (T2 LEFT JOIN T3 ON P2(T2,T3))
             ON P1(T1,T2)
  WHERE P(T1,T2,T3)
```

For this query, modify the nested-loop pattern to obtain:

```
FOR each row t1 in T1 {
  BOOL f1:=FALSE;
  FOR each row t2 in T2 such that P1(t1,t2) {
    BOOL f2:=FALSE;
    FOR each row t3 in T3 such that P2(t2,t3) {
      IF P(t1,t2,t3) {
        t:=t1||t2||t3; OUTPUT t;
      }
      f2=TRUE;
      f1=TRUE;
    }
    IF (!f2) {
      IF P(t1,t2,NULL) {
        t:=t1||t2||NULL; OUTPUT t;
      }
      f1=TRUE;
    }
  }
  IF (!f1) {
    IF P(t1,NULL,NULL) {
      t:=t1||NULL||NULL; OUTPUT t;
    }
  }
}
```

In general, for any nested loop for the first inner table in an outer join operation, a flag is introduced that is turned off before the loop and is checked after the loop. The flag is turned on when for the current row from the outer table a match from the table representing the inner operand is found. If at the end of the loop cycle the flag is still off, no match has been found for the current row of the outer table. In this case, the row is complemented by `NULL` values for the columns of the inner tables. The result row is passed to the final check for the output or into the next nested loop, but only if the row satisfies the join condition of all embedded outer joins.

In the example, the outer join table expressed by the following expression is embedded:

```
(T2 LEFT JOIN T3 ON P2(T2,T3))
```

For the query with inner joins, the optimizer could choose a different order of nested loops, such as this one:

```
FOR each row t3 in T3 {
  FOR each row t2 in T2 such that P2(t2,t3) {
    FOR each row t1 in T1 such that P1(t1,t2) {
      IF P(t1,t2,t3) {
         t:=t1||t2||t3; OUTPUT t;
      }
    }
  }
}
```

For queries with outer joins, the optimizer can choose only such an order where loops for outer tables precede loops for inner tables. Thus, for our query with outer joins, only one nesting order is possible. For the following query, the optimizer evaluates two different nestings. In both nestings, `T1` must be processed in the outer loop because it is used in an outer join. `T2` and `T3` are used in an inner join, so that join must be processed in the inner loop. However, because the join is an inner join, `T2` and `T3` can be processed in either order.

```
SELECT * T1 LEFT JOIN (T2,T3) ON P1(T1,T2) AND P2(T1,T3)
  WHERE P(T1,T2,T3)
```

One nesting evaluates `T2`, then `T3`:

```
FOR each row t1 in T1 {
  BOOL f1:=FALSE;
  FOR each row t2 in T2 such that P1(t1,t2) {
    FOR each row t3 in T3 such that P2(t1,t3) {
      IF P(t1,t2,t3) {
        t:=t1||t2||t3; OUTPUT t;
      }
      f1:=TRUE
    }
  }
  IF (!f1) {
    IF P(t1,NULL,NULL) {
      t:=t1||NULL||NULL; OUTPUT t;
    }
  }
}
```

The other nesting evaluates `T3`, then `T2`:

```
FOR each row t1 in T1 {
  BOOL f1:=FALSE;
  FOR each row t3 in T3 such that P2(t1,t3) {
    FOR each row t2 in T2 such that P1(t1,t2) {
      IF P(t1,t2,t3) {
        t:=t1||t2||t3; OUTPUT t;
      }
      f1:=TRUE
    }
  }
  IF (!f1) {
    IF P(t1,NULL,NULL) {
      t:=t1||NULL||NULL; OUTPUT t;
    }
  }
}
```

When discussing the nested-loop algorithm for inner joins, we omitted some details whose impact on the performance of query execution may be huge. We did not mention so-called “pushed-down” conditions. Suppose that our `WHERE` condition `P(T1,T2,T3)` can be represented by a conjunctive formula:

```
P(T1,T2,T2) = C1(T1) AND C2(T2) AND C3(T3).
```

In this case, MySQL actually uses the following nested-loop algorithm for the execution of the query with inner joins:

```
FOR each row t1 in T1 such that C1(t1) {
  FOR each row t2 in T2 such that P1(t1,t2) AND C2(t2)  {
    FOR each row t3 in T3 such that P2(t2,t3) AND C3(t3) {
      IF P(t1,t2,t3) {
         t:=t1||t2||t3; OUTPUT t;
      }
    }
  }
}
```

You see that each of the conjuncts `C1(T1)`, `C2(T2)`, `C3(T3)` are pushed out of the most inner loop to the most outer loop where it can be evaluated. If `C1(T1)` is a very restrictive condition, this condition pushdown may greatly reduce the number of rows from table `T1` passed to the inner loops. As a result, the execution time for the query may improve immensely.

For a query with outer joins, the `WHERE` condition is to be checked only after it has been found that the current row from the outer table has a match in the inner tables. Thus, the optimization of pushing conditions out of the inner nested loops cannot be applied directly to queries with outer joins. Here we must introduce conditional pushed-down predicates guarded by the flags that are turned on when a match has been encountered.

Recall this example with outer joins:

```
P(T1,T2,T3)=C1(T1) AND C(T2) AND C3(T3)
```

For that example, the nested-loop algorithm using guarded pushed-down conditions looks like this:

```
FOR each row t1 in T1 such that C1(t1) {
  BOOL f1:=FALSE;
  FOR each row t2 in T2
      such that P1(t1,t2) AND (f1?C2(t2):TRUE) {
    BOOL f2:=FALSE;
    FOR each row t3 in T3
        such that P2(t2,t3) AND (f1&&f2?C3(t3):TRUE) {
      IF (f1&&f2?TRUE:(C2(t2) AND C3(t3))) {
        t:=t1||t2||t3; OUTPUT t;
      }
      f2=TRUE;
      f1=TRUE;
    }
    IF (!f2) {
      IF (f1?TRUE:C2(t2) && P(t1,t2,NULL)) {
        t:=t1||t2||NULL; OUTPUT t;
      }
      f1=TRUE;
    }
  }
  IF (!f1 && P(t1,NULL,NULL)) {
      t:=t1||NULL||NULL; OUTPUT t;
  }
}
```

In general, pushed-down predicates can be extracted from join conditions such as `P1(T1,T2)` and `P(T2,T3)`. In this case, a pushed-down predicate is guarded also by a flag that prevents checking the predicate for the `NULL`-complemented row generated by the corresponding outer join operation.

Access by key from one inner table to another in the same nested join is prohibited if it is induced by a predicate from the `WHERE` condition.

8.2.1.8 Left Join and Right Join Optimization

MySQL implements an *`A LEFT JOIN B join_condition`* as follows:

- Table *B* is set to depend on table *A* and all tables on which *A* depends.
- Table *A* is set to depend on all tables (except *B*) that are used in the `LEFT JOIN` condition.
- The `LEFT JOIN` condition is used to decide how to retrieve rows from table *B*. (In other words, any condition in the `WHERE` clause is not used.)
- All standard join optimizations are performed, with the exception that a table is always read after all tables on which it depends. If there is a circular dependency, an error occurs.
- All standard `WHERE` optimizations are performed.
- If there is a row in *A* that matches the `WHERE` clause, but there is no row in *B* that matches the `ON` condition, an extra *B* row is generated with all columns set to `NULL`.
- If you use `LEFT JOIN` to find rows that do not exist in some table and you have the following test: *`col_name`* `IS NULL` in the `WHERE` part, where *`col_name`* is a column that is declared as `NOT NULL`, MySQL stops searching for more rows (for a particular key combination) after it has found one row that matches the `LEFT JOIN` condition.

The `RIGHT JOIN` implementation is analogous to that of `LEFT JOIN` with the table roles reversed. Right joins are converted to equivalent left joins, as described in Section 8.2.1.9, “Outer Join Simplification”.

The join optimizer calculates the order in which to join tables. The table read order forced by `LEFT JOIN` or `STRAIGHT_JOIN` helps the join optimizer do its work much more quickly, because there are fewer table permutations to check. This means that if you execute a query of the following type, MySQL does a full scan on `b` because the `LEFT JOIN` forces it to be read before `d`:

```
SELECT *
  FROM a JOIN b LEFT JOIN c ON (c.key=a.key)
  LEFT JOIN d ON (d.key=a.key)
  WHERE b.key=d.key;
```

The fix in this case is reverse the order in which `a` and `b` are listed in the `FROM` clause:

```
SELECT *
  FROM b JOIN a LEFT JOIN c ON (c.key=a.key)
  LEFT JOIN d ON (d.key=a.key)
  WHERE b.key=d.key;
```

For a `LEFT JOIN`, if the `WHERE` condition is always false for the generated `NULL` row, the `LEFT JOIN` is changed to a normal join. For example, the `WHERE` clause would be false in the following query if `t2.column1` were `NULL`:

```
SELECT * FROM t1 LEFT JOIN t2 ON (column1) WHERE t2.column2=5;
```

Therefore, it is safe to convert the query to a normal join:

```
SELECT * FROM t1, t2 WHERE t2.column2=5 AND t1.column1=t2.column1;
```

Now the optimizer can use table `t2` before table `t1` if doing so would result in a better query plan. To provide a hint about the table join order, use `STRAIGHT_JOIN`; see Section 13.2.10, “SELECT Syntax”. However, `STRAIGHT_JOIN` may prevent indexes from being used because it disables semi-join transformations; see Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”.

8.2.1.9 Outer Join Simplification

Table expressions in the `FROM` clause of a query are simplified in many cases.

At the parser stage, queries with right outer join operations are converted to equivalent queries containing only left join operations. In the general case, the conversion is performed such that this right join:

```
(T1, ...) RIGHT JOIN (T2, ...) ON P(T1, ..., T2, ...)
```

Becomes this equivalent left join:

```
(T2, ...) LEFT JOIN (T1, ...) ON P(T1, ..., T2, ...)
```

All inner join expressions of the form `T1 INNER JOIN T2 ON P(T1,T2)` are replaced by the list `T1,T2`, `P(T1,T2)` being joined as a conjunct to the `WHERE` condition (or to the join condition of the embedding join, if there is any).

When the optimizer evaluates plans for outer join operations, it takes into consideration only plans where, for each such operation, the outer tables are accessed before the inner tables. The optimizer choices are limited because only such plans enable outer joins to be executed using the nested-loop algorithm.

Consider a query of this form, where `R(T2)` greatly narrows the number of matching rows from table `T2`:

```
SELECT * T1 LEFT JOIN T2 ON P1(T1,T2)
  WHERE P(T1,T2) AND R(T2)
```

If the query is executed as written, the optimizer has no choice but to access the less-restricted table `T1` before the more-restricted table `T2`, which may produce a very inefficient execution plan.

Instead, MySQL converts the query to a query with no outer join operation if the `WHERE` condition is null-rejected. (That is, it converts the outer join to an inner join.) A condition is said to be null-rejected for an outer join operation if it evaluates to `FALSE` or `UNKNOWN` for any `NULL`-complemented row generated for the operation.

Thus, for this outer join:

```
T1 LEFT JOIN T2 ON T1.A=T2.A
```

Conditions such as these are null-rejected because they cannot be true for any `NULL`-complemented row (with `T2` columns set to `NULL`):

```
T2.B IS NOT NULL
T2.B > 3
T2.C <= T1.C
T2.B < 2 OR T2.C > 1
```

Conditions such as these are not null-rejected because they might be true for a `NULL`-complemented row:

```
T2.B IS NULL
T1.B < 3 OR T2.B IS NOT NULL
T1.B < 3 OR T2.B > 3
```

The general rules for checking whether a condition is null-rejected for an outer join operation are simple:

- It is of the form `A IS NOT NULL`, where `A` is an attribute of any of the inner tables
- It is a predicate containing a reference to an inner table that evaluates to `UNKNOWN` when one of its arguments is `NULL`

- It is a conjunction containing a null-rejected condition as a conjunct
- It is a disjunction of null-rejected conditions

A condition can be null-rejected for one outer join operation in a query and not null-rejected for another. In this query, the `WHERE` condition is null-rejected for the second outer join operation but is not null-rejected for the first one:

```
SELECT * FROM T1 LEFT JOIN T2 ON T2.A=T1.A
                 LEFT JOIN T3 ON T3.B=T1.B
  WHERE T3.C > 0
```

If the `WHERE` condition is null-rejected for an outer join operation in a query, the outer join operation is replaced by an inner join operation.

For example, in the preceding query, the second outer join is null-rejected and can be replaced by an inner join:

```
SELECT * FROM T1 LEFT JOIN T2 ON T2.A=T1.A
                 INNER JOIN T3 ON T3.B=T1.B
  WHERE T3.C > 0
```

For the original query, the optimizer evaluates only plans compatible with the single table-access order `T1,T2,T3`. For the rewritten query, it additionally considers the access order `T3,T1,T2`.

A conversion of one outer join operation may trigger a conversion of another. Thus, the query:

```
SELECT * FROM T1 LEFT JOIN T2 ON T2.A=T1.A
                 LEFT JOIN T3 ON T3.B=T2.B
  WHERE T3.C > 0
```

Is first converted to the query:

```
SELECT * FROM T1 LEFT JOIN T2 ON T2.A=T1.A
                 INNER JOIN T3 ON T3.B=T2.B
  WHERE T3.C > 0
```

Which is equivalent to the query:

```
SELECT * FROM (T1 LEFT JOIN T2 ON T2.A=T1.A), T3
  WHERE T3.C > 0 AND T3.B=T2.B
```

The remaining outer join operation can also be replaced by an inner join because the condition `T3.B=T2.B` is null-rejected. This results in a query with no outer joins at all:

```
SELECT * FROM (T1 INNER JOIN T2 ON T2.A=T1.A), T3
  WHERE T3.C > 0 AND T3.B=T2.B
```

Sometimes the optimizer succeeds in replacing an embedded outer join operation, but cannot convert the embedding outer join. The following query:

```
SELECT * FROM T1 LEFT JOIN
              (T2 LEFT JOIN T3 ON T3.B=T2.B)
              ON T2.A=T1.A
  WHERE T3.C > 0
```

Is converted to:

```
SELECT * FROM T1 LEFT JOIN
              (T2 INNER JOIN T3 ON T3.B=T2.B)
              ON T2.A=T1.A
```

```
  WHERE T3.C > 0
```

That can be rewritten only to the form still containing the embedding outer join operation:

```
SELECT * FROM T1 LEFT JOIN
             (T2,T3)
             ON (T2.A=T1.A AND T3.B=T2.B)
  WHERE T3.C > 0
```

Any attempt to convert an embedded outer join operation in a query must take into account the join condition for the embedding outer join together with the `WHERE` condition. In this query, the `WHERE` condition is not null-rejected for the embedded outer join, but the join condition of the embedding outer join `T2.A=T1.A AND T3.C=T1.C` is null-rejected:

```
SELECT * FROM T1 LEFT JOIN
             (T2 LEFT JOIN T3 ON T3.B=T2.B)
             ON T2.A=T1.A AND T3.C=T1.C
  WHERE T3.D > 0 OR T1.D > 0
```

Consequently, the query can be converted to:

```
SELECT * FROM T1 LEFT JOIN
             (T2, T3)
             ON T2.A=T1.A AND T3.C=T1.C AND T3.B=T2.B
  WHERE T3.D > 0 OR T1.D > 0
```

8.2.1.10 Multi-Range Read Optimization

Reading rows using a range scan on a secondary index can result in many random disk accesses to the base table when the table is large and not stored in the storage engine's cache. With the Disk-Sweep Multi-Range Read (MRR) optimization, MySQL tries to reduce the number of random disk access for range scans by first scanning the index only and collecting the keys for the relevant rows. Then the keys are sorted and finally the rows are retrieved from the base table using the order of the primary key. The motivation for Disk-sweep MRR is to reduce the number of random disk accesses and instead achieve a more sequential scan of the base table data.

The Multi-Range Read optimization provides these benefits:

- MRR enables data rows to be accessed sequentially rather than in random order, based on index tuples. The server obtains a set of index tuples that satisfy the query conditions, sorts them according to data row ID order, and uses the sorted tuples to retrieve data rows in order. This makes data access more efficient and less expensive.
- MRR enables batch processing of requests for key access for operations that require access to data rows through index tuples, such as range index scans and equi-joins that use an index for the join attribute. MRR iterates over a sequence of index ranges to obtain qualifying index tuples. As these results accumulate, they are used to access the corresponding data rows. It is not necessary to acquire all index tuples before starting to read data rows.

The MRR optimization is not supported with secondary indexes created on virtual generated columns. `InnoDB` supports secondary indexes on virtual generated columns.

The following scenarios illustrate when MRR optimization can be advantageous:

Scenario A: MRR can be used for `InnoDB` and `MyISAM` tables for index range scans and equi-join operations.

1. A portion of the index tuples are accumulated in a buffer.
2. The tuples in the buffer are sorted by their data row ID.
3. Data rows are accessed according to the sorted index tuple sequence.

Scenario B: MRR can be used for NDB tables for multiple-range index scans or when performing an equi-join by an attribute.

1. A portion of ranges, possibly single-key ranges, is accumulated in a buffer on the central node where the query is submitted.
2. The ranges are sent to the execution nodes that access data rows.
3. The accessed rows are packed into packages and sent back to the central node.
4. The received packages with data rows are placed in a buffer.
5. Data rows are read from the buffer.

When MRR is used, the Extra column in EXPLAIN output shows Using MRR.

InnoDB and MyISAM do not use MRR if full table rows need not be accessed to produce the query result. This is the case if results can be produced entirely on the basis on information in the index tuples (through a covering index); MRR provides no benefit.

Example query for which MRR can be used, assuming that there is an index on *(key_part1, key_part2)*:

```
SELECT * FROM t
  WHERE key_part1 >= 1000 AND key_part1 < 2000
  AND key_part2 = 10000;
```

The index consists of tuples of *(key_part1, key_part2)* values, ordered first by *key_part1* and then by *key_part2*.

Without MRR, an index scan covers all index tuples for the *key_part1* range from 1000 up to 2000, regardless of the *key_part2* value in these tuples. The scan does extra work to the extent that tuples in the range contain *key_part2* values other than 10000.

With MRR, the scan is broken up into multiple ranges, each for a single value of *key_part1* (1000, 1001, ... , 1999). Each of these scans need look only for tuples with *key_part2* = 10000. If the index contains many tuples for which *key_part2* is not 10000, MRR results in many fewer index tuples being read.

To express this using interval notation, the non-MRR scan must examine the index range [{1000, 10000}, {2000, MIN_INT}), which may include many tuples other than those for which *key_part2* = 10000. The MRR scan examines multiple single-point intervals [{1000, 10000}], ..., [{1999, 10000}], which includes only tuples with *key_part2* = 10000.

Two optimizer_switch system variable flags provide an interface to the use of MRR optimization. The mrr flag controls whether MRR is enabled. If mrr is enabled (on), the mrr_cost_based flag controls whether the optimizer attempts to make a cost-based choice between using and not using MRR (on) or uses MRR whenever possible (off). By default, mrr is on and mrr_cost_based is on. See Section 8.9.3, “Switchable Optimizations”.

For MRR, a storage engine uses the value of the read_rnd_buffer_size system variable as a guideline for how much memory it can allocate for its buffer. The engine uses up to read_rnd_buffer_size bytes and determines the number of ranges to process in a single pass.

8.2.1.11 Block Nested-Loop and Batched Key Access Joins

In MySQL, a Batched Key Access (BKA) Join algorithm is available that uses both index access to the joined table and a join buffer. The BKA algorithm supports inner join, outer join, and semi-join operations, including nested outer joins. Benefits of BKA include improved join performance due to more efficient table scanning. Also, the Block Nested-Loop (BNL) Join algorithm previously used only for inner joins is extended and can be employed for outer join and semi-join operations, including nested outer joins.

The following sections discuss the join buffer management that underlies the extension of the original BNL algorithm, the extended BNL algorithm, and the BKA algorithm. For information about semi-join strategies, see Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”

- Join Buffer Management for Block Nested-Loop and Batched Key Access Algorithms
- Block Nested-Loop Algorithm for Outer Joins and Semi-Joins
- Batched Key Access Joins
- Optimizer Hints for Block Nested-Loop and Batched Key Access Algorithms

Join Buffer Management for Block Nested-Loop and Batched Key Access Algorithms

MySQL can employ join buffers to execute not only inner joins without index access to the inner table, but also outer joins and semi-joins that appear after subquery flattening. Moreover, a join buffer can be effectively used when there is an index access to the inner table.

The join buffer management code slightly more efficiently utilizes join buffer space when storing the values of the interesting row columns: No additional bytes are allocated in buffers for a row column if its value is `NULL`, and the minimum number of bytes is allocated for any value of the `VARCHAR` type.

The code supports two types of buffers, regular and incremental. Suppose that join buffer `B1` is employed to join tables `t1` and `t2` and the result of this operation is joined with table `t3` using join buffer `B2`:

- A regular join buffer contains columns from each join operand. If `B2` is a regular join buffer, each row *`r`* put into `B2` is composed of the columns of a row *`r1`* from `B1` and the interesting columns of a matching row *`r2`* from table `t3`.
- An incremental join buffer contains only columns from rows of the table produced by the second join operand. That is, it is incremental to a row from the first operand buffer. If `B2` is an incremental join buffer, it contains the interesting columns of the row *`r2`* together with a link to the row *`r1`* from `B1`.

Incremental join buffers are always incremental relative to a join buffer from an earlier join operation, so the buffer from the first join operation is always a regular buffer. In the example just given, the buffer `B1` used to join tables `t1` and `t2` must be a regular buffer.

Each row of the incremental buffer used for a join operation contains only the interesting columns of a row from the table to be joined. These columns are augmented with a reference to the interesting columns of the matched row from the table produced by the first join operand. Several rows in the incremental buffer can refer to the same row *`r`* whose columns are stored in the previous join buffers insofar as all these rows match row *`r`*.

Incremental buffers enable less frequent copying of columns from buffers used for previous join operations. This provides a savings in buffer space because in the general case a row produced by the first join operand can be matched by several rows produced by the second join operand. It is unnecessary to make several copies of a row from the first operand. Incremental buffers also provide a savings in processing time due to the reduction in copying time.

The `block_nested_loop` and `batched_key_access` flags of the `optimizer_switch` system variable control how the optimizer uses the Block Nested-Loop and Batched Key Access join algorithms. By default, `block_nested_loop` is `on` and `batched_key_access` is `off`. See Section 8.9.3, “Switchable Optimizations”. Optimizer hints may also be applied; see Optimizer Hints for Block Nested-Loop and Batched Key Access Algorithms.

For information about semi-join strategies, see Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”

Block Nested-Loop Algorithm for Outer Joins and Semi-Joins

The original implementation of the MySQL BNL algorithm is extended to support outer join and semi-join operations.

When these operations are executed with a join buffer, each row put into the buffer is supplied with a match flag.

If an outer join operation is executed using a join buffer, each row of the table produced by the second operand is checked for a match against each row in the join buffer. When a match is found, a new extended row is formed (the original row plus columns from the second operand) and sent for further extensions by the remaining join operations. In addition, the match flag of the matched row in the buffer is enabled. After all rows of the table to be joined have been examined, the join buffer is scanned. Each row from the buffer that does not have its match flag enabled is extended by `NULL` complements (`NULL` values for each column in the second operand) and sent for further extensions by the remaining join operations.

The `block_nested_loop` flag of the `optimizer_switch` system variable controls how the optimizer uses the Block Nested-Loop algorithm. By default, `block_nested_loop` is `on`. See Section 8.9.3, “Switchable Optimizations”. Optimizer hints may also be applied; see Optimizer Hints for Block Nested-Loop and Batched Key Access Algorithms.

In `EXPLAIN` output, use of BNL for a table is signified when the `Extra` value contains `Using join buffer (Block Nested Loop)` and the `type` value is `ALL`, `index`, or `range`.

For information about semi-join strategies, see Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”

Batched Key Access Joins

MySQL implements a method of joining tables called the Batched Key Access (BKA) join algorithm. BKA can be applied when there is an index access to the table produced by the second join operand. Like the BNL join algorithm, the BKA join algorithm employs a join buffer to accumulate the interesting columns of the rows produced by the first operand of the join operation. Then the BKA algorithm builds keys to access the table to be joined for all rows in the buffer and submits these keys in a batch to the database engine for index lookups. The keys are submitted to the engine through the Multi-Range Read (MRR) interface (see Section 8.2.1.10, “Multi-Range Read Optimization”). After submission of the keys, the MRR engine functions perform lookups in the index in an optimal way, fetching the rows of the joined table found by these keys, and starts feeding the BKA join algorithm with matching rows. Each matching row is coupled with a reference to a row in the join buffer.

When BKA is used, the value of `join_buffer_size` defines how large the batch of keys is in each request to the storage engine. The larger the buffer, the more sequential access will be to the right hand table of a join operation, which can significantly improve performance.

For BKA to be used, the `batched_key_access` flag of the `optimizer_switch` system variable must be set to `on`. BKA uses MRR, so the `mrr` flag must also be `on`. Currently, the cost estimation for MRR is too pessimistic. Hence, it is also necessary for `mrr_cost_based` to be `off` for BKA to be used. The following setting enables BKA:

```
mysql> SET optimizer_switch='mrr=on,mrr_cost_based=off,batched_key_access=on';
```

There are two scenarios by which MRR functions execute:

- The first scenario is used for conventional disk-based storage engines such as `InnoDB` and `MyISAM`. For these engines, usually the keys for all rows from the join buffer are submitted to the MRR interface at once. Engine-specific MRR functions perform index lookups for the submitted keys, get row IDs (or primary keys) from them, and then fetch rows for all these selected row IDs one by one by request from BKA algorithm. Every row is returned with an association reference that enables access to the matched row in the join buffer. The rows are fetched by the MRR functions in an optimal way: They are fetched in the row ID (primary key) order. This improves performance because reads are in disk order rather than random order.

- The second scenario is used for remote storage engines such as NDB. A package of keys for a portion of rows from the join buffer, together with their associations, is sent by a MySQL Server (SQL node) to MySQL Cluster data nodes. In return, the SQL node receives a package (or several packages) of matching rows coupled with corresponding associations. The BKA join algorithm takes these rows and builds new joined rows. Then a new set of keys is sent to the data nodes and the rows from the returned packages are used to build new joined rows. The process continues until the last keys from the join buffer are sent to the data nodes, and the SQL node has received and joined all rows matching these keys. This improves performance because fewer key-bearing packages sent by the SQL node to the data nodes means fewer round trips between it and the data nodes to perform the join operation.

With the first scenario, a portion of the join buffer is reserved to store row IDs (primary keys) selected by index lookups and passed as a parameter to the MRR functions.

There is no special buffer to store keys built for rows from the join buffer. Instead, a function that builds the key for the next row in the buffer is passed as a parameter to the MRR functions.

In EXPLAIN output, use of BKA for a table is signified when the Extra value contains Using join buffer (Batched Key Access) and the type value is ref or eq_ref.

Optimizer Hints for Block Nested-Loop and Batched Key Access Algorithms

In addition to using the optimizer_switch system variable to control optimizer use of the BNL and BKA algorithms session-wide, MySQL supports optimizer hints to influence the optimizer on a per-statement basis. See Section 8.9.2, “Optimizer Hints”.

To use a BNL or BKA hint to enable join buffering for any inner table of an outer join, join buffering must be enabled for all inner tables of the outer join.

8.2.1.12 IS NULL Optimization

MySQL can perform the same optimization on *col_name* IS NULL that it can use for *col_name* = *constant_value*. For example, MySQL can use indexes and ranges to search for NULL with IS NULL.

Examples:

```
SELECT * FROM tbl_name WHERE key_col IS NULL;

SELECT * FROM tbl_name WHERE key_col <=> NULL;

SELECT * FROM tbl_name
  WHERE key_col=const1 OR key_col=const2 OR key_col IS NULL;
```

If a WHERE clause includes a *col_name* IS NULL condition for a column that is declared as NOT NULL, that expression is optimized away. This optimization does not occur in cases when the column might produce NULL anyway; for example, if it comes from a table on the right side of a LEFT JOIN.

MySQL can also optimize the combination *col_name* = *expr* OR *col_name* IS NULL, a form that is common in resolved subqueries. EXPLAIN shows ref_or_null when this optimization is used.

This optimization can handle one IS NULL for any key part.

Some examples of queries that are optimized, assuming that there is an index on columns a and b of table t2:

```
SELECT * FROM t1 WHERE t1.a=expr OR t1.a IS NULL;

SELECT * FROM t1, t2 WHERE t1.a=t2.a OR t2.a IS NULL;

SELECT * FROM t1, t2
  WHERE (t1.a=t2.a OR t2.a IS NULL) AND t2.b=t1.b;
```

```
SELECT * FROM t1, t2
  WHERE t1.a=t2.a AND (t2.b=t1.b OR t2.b IS NULL);

SELECT * FROM t1, t2
  WHERE (t1.a=t2.a AND t2.a IS NULL AND ...)
  OR (t1.a=t2.a AND t2.a IS NULL AND ...);
```

`ref_or_null` works by first doing a read on the reference key, and then a separate search for rows with a `NULL` key value.

The optimization can handle only one `IS NULL` level. In the following query, MySQL uses key lookups only on the expression `(t1.a=t2.a AND t2.a IS NULL)` and is not able to use the key part on `b`:

```
SELECT * FROM t1, t2
  WHERE (t1.a=t2.a AND t2.a IS NULL)
  OR (t1.b=t2.b AND t2.b IS NULL);
```

8.2.1.13 ORDER BY Optimization

This section describes when MySQL can use an index to satisfy an `ORDER BY` clause, the `filesort` algorithms used when an index cannot be used, and execution plan information available from the optimizer about `ORDER BY`.

- Use of Indexes to Satisfy ORDER BY
- Optimization Using filesort
- The Original filesort Algorithm
- The Modified filesort Algorithm
- The In-Memory filesort Algorithm
- Comparison of filesort Algorithms
- Influencing ORDER BY Optimization
- ORDER BY Execution Plan Information Available

Use of Indexes to Satisfy ORDER BY

In some cases, MySQL can use an index to satisfy an `ORDER BY` clause without doing extra sorting.

The index can also be used even if the `ORDER BY` does not match the index exactly, as long as all unused portions of the index and all extra `ORDER BY` columns are constants in the `WHERE` clause. The following queries use the index to resolve the `ORDER BY` part:

```
SELECT * FROM t1
  ORDER BY key_part1, key_part2;

SELECT * FROM t1
  WHERE key_part1 = constant
  ORDER BY key_part2;

SELECT * FROM t1
  ORDER BY key_part1 DESC, key_part2 DESC;

SELECT * FROM t1
  WHERE key_part1 = 1
  ORDER BY key_part1 DESC, key_part2 DESC;

SELECT * FROM t1
  WHERE key_part1 > constant
  ORDER BY key_part1 ASC;
```

```
SELECT * FROM t1
  WHERE key_part1 < constant
  ORDER BY key_part1 DESC;

SELECT * FROM t1
  WHERE key_part1 = constant1 AND key_part2 > constant2
  ORDER BY key_part2;
```

If the query mixes `ASC` and `DESC`, the optimizer can use an index on the columns if the index also uses corresponding mixed ascending and descending columns:

```
SELECT * FROM t1 ORDER BY key_part1 DESC, key_part2 ASC;
```

The optimizer can use an index on (*`key_part1`*, *`key_part2`*) if *`key_part1`* is descending and *`key_part2`* is ascending. It can also use an index on those columns (with a backward scan) if *`key_part1`* is ascending and *`key_part2`* is descending. See Section 8.3.12, “Descending Indexes”

In some cases, MySQL *cannot* use indexes to resolve the `ORDER BY`, although it may still use indexes to find the rows that match the `WHERE` clause. Examples:

- The query uses `ORDER BY` on different indexes:

  ```
  SELECT * FROM t1 ORDER BY key1, key2;
  ```

- The query uses `ORDER BY` on nonconsecutive parts of an index:

  ```
  SELECT * FROM t1 WHERE key2=constant ORDER BY key_part1, key_part3;
  ```

- The index used to fetch the rows differs from the one used in the `ORDER BY`:

  ```
  SELECT * FROM t1 WHERE key2=constant ORDER BY key1;
  ```

- The query uses `ORDER BY` with an expression that includes terms other than the index column name:

  ```
  SELECT * FROM t1 ORDER BY ABS(key);
  SELECT * FROM t1 ORDER BY -key;
  ```

- The query joins many tables, and the columns in the `ORDER BY` are not all from the first nonconstant table that is used to retrieve rows. (This is the first table in the `EXPLAIN` output that does not have a `const` join type.)

- The query has different `ORDER BY` and `GROUP BY` expressions.

- There is an index on only a prefix of a column named in the `ORDER BY` clause. In this case, the index cannot be used to fully resolve the sort order. For example, if only the first 10 bytes of a `CHAR(20)` column are indexed, the index cannot distinguish values past the 10th byte and a `filesort` is needed.

- The index does not store rows in order. For example, this is true for a `HASH` index in a `MEMORY` table.

Availability of an index for sorting may be affected by the use of column aliases. Suppose that the column `t1.a` is indexed. In this statement, the name of the column in the select list is `a`. It refers to `t1.a`, as does the reference to `a` in the `ORDER BY`, so the index on `t1.a` can be used:

```
SELECT a FROM t1 ORDER BY a;
```

In this statement, the name of the column in the select list is also `a`, but it is the alias name. It refers to `ABS(a)`, as does the reference to `a` in the `ORDER BY`, so the index on `t1.a` cannot be used:

```
```

```
SELECT ABS(a) AS a FROM t1 ORDER BY a;
```

In the following statement, the `ORDER BY` refers to a name that is not the name of a column in the select list. But there is a column in `t1` named `a`, so the `ORDER BY` refers to `t1.a` and the index on `t1.a` can be used. (The resulting sort order may be completely different from the order for `ABS(a)`, of course.)

```
SELECT ABS(a) AS b FROM t1 ORDER BY a;
```

By default, MySQL sorts all `GROUP BY col1, col2, ...` queries as if you specified `ORDER BY col1, col2, ...` in the query as well. If you include an explicit `ORDER BY` clause that contains the same column list, MySQL optimizes it away without any speed penalty, although the sorting still occurs.

If a query includes `GROUP BY` but you want to avoid the overhead of sorting the result, you can suppress sorting by specifying `ORDER BY NULL`. For example:

```
INSERT INTO foo
SELECT a, COUNT(*) FROM bar GROUP BY a ORDER BY NULL;
```

The optimizer may still choose to use sorting to implement grouping operations. `ORDER BY NULL` suppresses sorting of the result, not prior sorting done by grouping operations to determine the result.

Note

Previously, relying on implicit `GROUP BY` sorting was deprecated but `GROUP BY` did sort by default (that is, in the absence of `ASC` or `DESC` designators). In MySQL 8.0, `GROUP BY` no longer sorts by default, so query results may differ from previous MySQL versions. To produce a given sort order, use explicit `ASC` or `DESC` designators for `GROUP BY` columns or provide an `ORDER BY` clause.

Optimization Using filesort

MySQL has multiple `filesort` algorithms for sorting and retrieving results. The original algorithm uses only the `ORDER BY` columns. The modified algorithm uses not just the `ORDER BY` columns, but all columns referenced by the query. There is also an algorithm for small result sets that sorts in memory using the sort buffer as a priority queue without a merge file.

The optimizer selects which `filesort` algorithm to use. It normally uses the modified algorithm except when `BLOB` or `TEXT` columns are involved, in which case it uses the original algorithm. For each algorithm, the sort buffer size is the `sort_buffer_size` system variable value.

The Original filesort Algorithm

The original `filesort` algorithm works as follows:

1. Read all rows according to key or by table scanning. Skip rows that do not match the `WHERE` clause.
2. For each row, store in the sort buffer a tuple consisting of a pair of values (the sort key value and the row ID).
3. If all pairs fit into the sort buffer, no temporary file is created. Otherwise, when the sort buffer becomes full, run a quicksort on it in memory and write it to a temporary file. Save a pointer to the sorted block.
4. Repeat the preceding steps until all rows have been read.
5. Do a multi-merge of up to `MERGEBUFF` (7) regions to one block in another temporary file. Repeat until all blocks from the first file are in the second file.
6. Repeat the following until there are fewer than `MERGEBUFF2` (15) blocks left.

7. On the last multi-merge, only the row ID (the last part of the value pair) is written to a result file.

8. Read the rows in sorted order using the row IDs in the result file. To optimize this, read in a large block of row IDs, sort them, and use them to read the rows in sorted order into a row buffer. The row buffer size is the `read_rnd_buffer_size` system variable value. The code for this step is in the `sql/records.cc` source file.

One problem with this approach is that it reads rows twice: Once during `WHERE` clause evaluation, and again after sorting the value pairs. And even if the rows were accessed successively the first time (for example, if a table scan is done), the second time they are accessed randomly. (The sort keys are ordered, but the row positions are not.)

The Modified filesort Algorithm

The modified `filesort` algorithm incorporates an optimization to avoid reading the rows twice: It records the sort key value, but instead of the row ID, it records the columns referenced by the query. The modified `filesort` algorithm works like this:

1. Read the rows that match the `WHERE` clause.

2. For each row, store in the sort buffer a tuple consisting of the sort key value and the columns referenced by the query.

3. When the sort buffer becomes full, sort the tuples by sort key value in memory and write it to a temporary file.

4. After merge-sorting the temporary file, retrieve the rows in sorted order, but read the columns required by the query directly from the sorted tuples rather than by accessing the table a second time.

The tuples used by the modified `filesort` algorithm are longer than the pairs used by the original algorithm, and fewer of them fit in the sort buffer. As a result, it is possible for the extra I/O to make the modified approach slower, not faster. To avoid a slowdown, the optimizer uses the modified algorithm only if the total size of the extra columns in the sort tuple does not exceed the value of the `max_length_for_sort_data` system variable. (A symptom of setting the value of this variable too high is a combination of high disk activity and low CPU activity.)

The modified `filesort` algorithm includes an additional optimization designed to enable more tuples to fit into the sort buffer: For additional columns of type `CHAR` or `VARCHAR`, or any nullable fixed-size data type, the values are packed. For example, without packing, a `VARCHAR(255)` column value containing only 3 characters takes 255 characters in the sort buffer. With packing, the value requires only 3 characters plus a two-byte length indicator. `NULL` values require only a bitmask.

For data containing packable strings shorter than the maximum column length or many `NULL` values, more records fit into the sort buffer. This improves in-memory sorting of the sort buffer and performance of disk-based temporary file merge sorting.

In edge cases, packing may be disadvantageous: If packable strings are the maximum column length or there are few `NULL` values, the space required for the length indicators reduces the number of records that fit into the sort buffer and sorting is slower in memory and on disk.

The In-Memory filesort Algorithm

For queries (and subqueries) of the following form, the optimizer can use `filesort` to efficiently handle the `ORDER BY` operation in memory without a merge file:

```
SELECT ... FROM single_table ... ORDER BY non_index_column [DESC] LIMIT [M,]N;
```

That type of query is common in web applications that display only a few rows from a larger result set. For example:

```
SELECT col1, ... FROM t1 ... ORDER BY name LIMIT 10;
SELECT col1, ... FROM t1 ... ORDER BY RAND() LIMIT 15;
```

The sort buffer has a size of `sort_buffer_size`. If the sort elements for *N* rows are small enough to fit in the sort buffer (*M*+*N* rows if *M* was specified), the server can avoid using a merge file and performs an in-memory sort by treating the sort buffer as a priority queue:

1. Scan the table, inserting the select list columns from each selected row in sorted order in the queue. If the queue is full, bump out the last row in the sort order.
2. Return the first *N* rows from the queue. (If *M* was specified, skip the first *M* rows and return the next *N* rows.)

Absent that optimization, the server performs this operation by using a merge file for the sort:

1. Scan the table, repeating these steps through the end of the table:
   - Select rows until the sort buffer is filled.
   - Write the first *N* rows in the buffer (*M*+*N* rows if *M* was specified) to a merge file.
2. Sort the merge file and return the first *N* rows. (If *M* was specified, skip the first *M* rows and return the next *N* rows.)

The cost of the table-scan operation is the same for the queue and merge-file methods, so the optimizer chooses between methods based on other costs:

- The queue method involves more CPU for inserting rows into the queue in order.
- The merge-file method has I/O costs to write and read the file and CPU cost to sort it.

The optimizer considers the balance between these factors for particular values of *M*, *N*, and the row size.

An `ORDER BY` with and without `LIMIT` may return rows in different orders, as discussed in Section 8.2.1.16, “LIMIT Query Optimization”.

Comparison of filesort Algorithms

Suppose that a table `t1` has four `VARCHAR` columns `a`, `b`, `c`, and `d` and that the optimizer uses `filesort` for this query:

```
SELECT * FROM t1 ORDER BY a, b;
```

The query sorts by `a` and `b`, but returns all columns, so the columns referenced by the query are `a`, `b`, `c`, and `d`. Depending on which `filesort` algorithm the optimizer chooses, the query executes as follows:

For the original algorithm, sort buffer tuples have these contents:

```
(fixed size a value, fixed size b value,
row ID into t1)
```

The optimizer sorts on the fixed size values. After sorting, the optimizer reads the tuples in order and uses the row ID in each tuple to read rows from `t1` to obtain the select list column values.

For the modified algorithm without packing, sort buffer tuples have these contents:

```
(fixed size a value, fixed size b value,
a value, b value, c value, d value)
```

The optimizer sorts on the fixed size values. After sorting, the optimizer reads the tuples in order and uses the values for `a`, `b`, `c`, and `d` to obtain the select list column values without reading `t1` again.

For the modified algorithm with packing, sort buffer tuples have these contents:

```
(fixed size a value, fixed size b value,
a length, packed a value, b length, packed b value,
c length, packed c value, d length, packed d value)
```

If any of `a`, `b`, `c`, or `d` are `NULL`, they take no space in the sort buffer other than in the bitmask.

The optimizer sorts on the fixed size values. After sorting, the optimizer reads the tuples in order and uses the values for `a`, `b`, `c`, and `d` to obtain the select list column values without reading `t1` again.

Influencing ORDER BY Optimization

For slow `ORDER BY` queries for which `filesort` is not used, try lowering `max_length_for_sort_data` to a value that is appropriate to trigger a `filesort`.

To increase `ORDER BY` speed, check whether you can get MySQL to use indexes rather than an extra sorting phase. If this is not possible, you can try the following strategies:

- Increase the `sort_buffer_size` variable value. Ideally, the value should be large enough for the entire result set to fit in the sort buffer (to avoid writes to disk and merge passes), but at minimum the value must be large enough to accommodate fifteen tuples.

  Take into account that the size of column values stored in the sort buffer is affected by the `max_sort_length` system variable value. For example, if tuples store values of long string columns and you increase the value of `max_sort_length`, the size of sort buffer tuples increases as well and may require you to increase `sort_buffer_size`. For column values calculated as a result of string expressions (such as those that invoke a string-valued function), the `filesort` algorithm cannot tell the maximum length of expression values, so it must allocate `max_sort_length` bytes for each tuple.

  To monitor the number of merge passes, check the `Sort_merge_passes` status variable.

- Increase the `read_rnd_buffer_size` variable value.

- Use less RAM per row by declaring columns only as large as they need to be to hold the values stored in them. For example, `CHAR(16)` is better than `CHAR(200)` if values never exceed 16 characters.

- Change the `tmpdir` system variable to point to a dedicated file system with large amounts of free space. The variable value can list several paths that are used in round-robin fashion; you can use this feature to spread the load across several directories. Separate the paths by colon characters (`:`) on Unix and semicolon characters (`;`) on Windows. The paths should name directories in file systems located on different *physical* disks, not different partitions on the same disk.

ORDER BY Execution Plan Information Available

With `EXPLAIN SELECT ... ORDER BY`, you can check whether MySQL can use indexes to resolve the query. It cannot if you see `Using filesort` in the `Extra` column. See Section 8.8.1, “Optimizing Queries with EXPLAIN”. Filesort uses a fixed-length row-storage format similar to that used by the `MEMORY` storage engine. Variable-length types such as `VARCHAR` are stored using a fixed length.

If a `filesort` is done, `EXPLAIN` output includes `Using filesort` in the `Extra` column. Also, optimizer trace output includes a `filesort_summary` block. For example:

```
"filesort_summary": {
  "rows": 100,
  "examined_rows": 100,
  "number_of_tmp_files": 0,
```

```
  "sort_buffer_size": 25192,
  "sort_mode": "<sort_key, packed_additional_fields>"
}
```

The `sort_mode` value provides information about the `filesort` algorithm used and the contents of tuples in the sort buffer:

- `<sort_key, rowid>`: This indicates use of the original algorithm. Sort buffer tuples are pairs that contain the sort key value and row ID of the original table row. Tuples are sorted by sort key value and the row ID is used to read the row from the table.
- `<sort_key, additional_fields>`: This indicates use of the modified algorithm. Sort buffer tuples contain the sort key value and columns referenced by the query. Tuples are sorted by sort key value and column values are read directly from the tuple.
- `<sort_key, packed_additional_fields>`: This indicates use of the modified algorithm. Sort buffer tuples contain the sort key value and packed columns referenced by the query. Tuples are sorted by sort key value and column values are read directly from the tuple.

`EXPLAIN` does not distinguish between use of the modified and in-memory (priority queue) algorithms. Use of the in-memory algorithm can be seen in optimizer trace output. Look for `filesort_priority_queue_optimization`.

For information about the optimizer trace, see MySQL Internals: Tracing the Optimizer.

8.2.1.14 GROUP BY Optimization

The most general way to satisfy a `GROUP BY` clause is to scan the whole table and create a new temporary table where all rows from each group are consecutive, and then use this temporary table to discover groups and apply aggregate functions (if any). In some cases, MySQL is able to do much better than that and to avoid creation of temporary tables by using index access.

The most important preconditions for using indexes for `GROUP BY` are that all `GROUP BY` columns reference attributes from the same index, and that the index stores its keys in order (for example, this is a `BTREE` index and not a `HASH` index). Whether use of temporary tables can be replaced by index access also depends on which parts of an index are used in a query, the conditions specified for these parts, and the selected aggregate functions.

There are two ways to execute a `GROUP BY` query through index access, as detailed in the following sections. In the first method, the grouping operation is applied together with all range predicates (if any). The second method first performs a range scan, and then groups the resulting tuples.

In MySQL, `GROUP BY` is used for sorting, so the server may also apply `ORDER BY` optimizations to grouping. However, relying on implicit `GROUP BY` sorting is deprecated. See Section 8.2.1.13, “ORDER BY Optimization”.

- Loose Index Scan
- Tight Index Scan

Loose Index Scan

The most efficient way to process `GROUP BY` is when an index is used to directly retrieve the grouping columns. With this access method, MySQL uses the property of some index types that the keys are ordered (for example, `BTREE`). This property enables use of lookup groups in an index without having to consider all keys in the index that satisfy all `WHERE` conditions. This access method considers only a fraction of the keys in an index, so it is called a *loose index scan*. When there is no `WHERE` clause, a loose index scan reads as many keys as the number of groups, which may be a much smaller number than that of all keys. If the `WHERE` clause contains range predicates (see the discussion of the `range` join type in Section 8.8.1, “Optimizing Queries with EXPLAIN”), a loose index scan looks up the first key of each group that satisfies the range conditions, and again reads the least possible number of keys. This is possible under the following conditions:

- The query is over a single table.
- The `GROUP BY` names only columns that form a leftmost prefix of the index and no other columns. (If, instead of `GROUP BY`, the query has a `DISTINCT` clause, all distinct attributes refer to columns that form a leftmost prefix of the index.) For example, if a table `t1` has an index on `(c1,c2,c3)`, loose index scan is applicable if the query has `GROUP BY c1, c2,`. It is not applicable if the query has `GROUP BY c2, c3` (the columns are not a leftmost prefix) or `GROUP BY c1, c2, c4` (`c4` is not in the index).
- The only aggregate functions used in the select list (if any) are `MIN()` and `MAX()`, and all of them refer to the same column. The column must be in the index and must immediately follow the columns in the `GROUP BY`.
- Any other parts of the index than those from the `GROUP BY` referenced in the query must be constants (that is, they must be referenced in equalities with constants), except for the argument of `MIN()` or `MAX()` functions.
- For columns in the index, full column values must be indexed, not just a prefix. For example, with `c1 VARCHAR(20), INDEX (c1(10))`, the index cannot be used for loose index scan.

If loose index scan is applicable to a query, the `EXPLAIN` output shows `Using index for group-by` in the `Extra` column.

Assume that there is an index `idx(c1,c2,c3)` on table `t1(c1,c2,c3,c4)`. The loose index scan access method can be used for the following queries:

```
SELECT c1, c2 FROM t1 GROUP BY c1, c2;
SELECT DISTINCT c1, c2 FROM t1;
SELECT c1, MIN(c2) FROM t1 GROUP BY c1;
SELECT c1, c2 FROM t1 WHERE c1 < const GROUP BY c1, c2;
SELECT MAX(c3), MIN(c3), c1, c2 FROM t1 WHERE c2 > const GROUP BY c1, c2;
SELECT c2 FROM t1 WHERE c1 < const GROUP BY c1, c2;
SELECT c1, c2 FROM t1 WHERE c3 = const GROUP BY c1, c2;
```

The following queries cannot be executed with this quick select method, for the reasons given:

- There are aggregate functions other than `MIN()` or `MAX()`:

  ```
  SELECT c1, SUM(c2) FROM t1 GROUP BY c1;
  ```

- The columns in the `GROUP BY` clause do not form a leftmost prefix of the index:

  ```
  SELECT c1, c2 FROM t1 GROUP BY c2, c3;
  ```

- The query refers to a part of a key that comes after the `GROUP BY` part, and for which there is no equality with a constant:

  ```
  SELECT c1, c3 FROM t1 GROUP BY c1, c2;
  ```

  Were the query to include `WHERE c3 = const`, loose index scan could be used.

The loose index scan access method can be applied to other forms of aggregate function references in the select list, in addition to the `MIN()` and `MAX()` references already supported:

- `AVG(DISTINCT)`, `SUM(DISTINCT)`, and `COUNT(DISTINCT)` are supported. `AVG(DISTINCT)` and `SUM(DISTINCT)` take a single argument. `COUNT(DISTINCT)` can have more than one column argument.
- There must be no `GROUP BY` or `DISTINCT` clause in the query.
- The loose scan limitations described earlier still apply.

Assume that there is an index `idx(c1,c2,c3)` on table `t1(c1,c2,c3,c4)`. The loose index scan access method can be used for the following queries:

```
SELECT COUNT(DISTINCT c1), SUM(DISTINCT c1) FROM t1;

SELECT COUNT(DISTINCT c1, c2), COUNT(DISTINCT c2, c1) FROM t1;
```

Tight Index Scan

A tight index scan may be either a full index scan or a range index scan, depending on the query conditions.

When the conditions for a loose index scan are not met, it still may be possible to avoid creation of temporary tables for `GROUP BY` queries. If there are range conditions in the `WHERE` clause, this method reads only the keys that satisfy these conditions. Otherwise, it performs an index scan. Because this method reads all keys in each range defined by the `WHERE` clause, or scans the whole index if there are no range conditions, we term it a *tight index scan*. With a tight index scan, the grouping operation is performed only after all keys that satisfy the range conditions have been found.

For this method to work, it is sufficient that there is a constant equality condition for all columns in a query referring to parts of the key coming before or in between parts of the `GROUP BY` key. The constants from the equality conditions fill in any “gaps” in the search keys so that it is possible to form complete prefixes of the index. These index prefixes then can be used for index lookups. If we require sorting of the `GROUP BY` result, and it is possible to form search keys that are prefixes of the index, MySQL also avoids extra sorting operations because searching with prefixes in an ordered index already retrieves all the keys in order.

Assume that there is an index `idx(c1,c2,c3)` on table `t1(c1,c2,c3,c4)`. The following queries do not work with the loose index scan access method described earlier, but still work with the tight index scan access method.

- There is a gap in the `GROUP BY`, but it is covered by the condition `c2 = 'a'`:

  ```
  SELECT c1, c2, c3 FROM t1 WHERE c2 = 'a' GROUP BY c1, c3;
  ```

- The `GROUP BY` does not begin with the first part of the key, but there is a condition that provides a constant for that part:

  ```
  SELECT c1, c2, c3 FROM t1 WHERE c1 = 'a' GROUP BY c2, c3;
  ```

8.2.1.15 DISTINCT Optimization

`DISTINCT` combined with `ORDER BY` needs a temporary table in many cases.

Because `DISTINCT` may use `GROUP BY`, learn how MySQL works with columns in `ORDER BY` or `HAVING` clauses that are not part of the selected columns. See Section 12.18.3, “MySQL Handling of GROUP BY”.

In most cases, a `DISTINCT` clause can be considered as a special case of `GROUP BY`. For example, the following two queries are equivalent:

```
SELECT DISTINCT c1, c2, c3 FROM t1
WHERE c1 > const;

SELECT c1, c2, c3 FROM t1
WHERE c1 > const GROUP BY c1, c2, c3;
```

Due to this equivalence, the optimizations applicable to `GROUP BY` queries can be also applied to queries with a `DISTINCT` clause. Thus, for more details on the optimization possibilities for `DISTINCT` queries, see Section 8.2.1.14, “GROUP BY Optimization”.

When combining `LIMIT row_count` with `DISTINCT`, MySQL stops as soon as it finds *`row_count`* unique rows.

If you do not use columns from all tables named in a query, MySQL stops scanning any unused tables as soon as it finds the first match. In the following case, assuming that `t1` is used before `t2` (which you can check with `EXPLAIN`), MySQL stops reading from `t2` (for any particular row in `t1`) when it finds the first row in `t2`:

```
SELECT DISTINCT t1.a FROM t1, t2 where t1.a=t2.a;
```

8.2.1.16 LIMIT Query Optimization

If you need only a specified number of rows from a result set, use a `LIMIT` clause in the query, rather than fetching the whole result set and throwing away the extra data.

MySQL sometimes optimizes a query that has a `LIMIT row_count` clause and no `HAVING` clause:

- If you select only a few rows with `LIMIT`, MySQL uses indexes in some cases when normally it would prefer to do a full table scan.
- If you combine `LIMIT row_count` with `ORDER BY`, MySQL stops sorting as soon as it has found the first *`row_count`* rows of the sorted result, rather than sorting the entire result. If ordering is done by using an index, this is very fast. If a filesort must be done, all rows that match the query without the `LIMIT` clause are selected, and most or all of them are sorted, before the first *`row_count`* are found. After the initial rows have been found, MySQL does not sort any remainder of the result set.

  One manifestation of this behavior is that an `ORDER BY` query with and without `LIMIT` may return rows in different order, as described later in this section.
- If you combine `LIMIT row_count` with `DISTINCT`, MySQL stops as soon as it finds *`row_count`* unique rows.
- In some cases, a `GROUP BY` can be resolved by reading the index in order (or doing a sort on the index), then calculating summaries until the index value changes. In this case, `LIMIT row_count` does not calculate any unnecessary `GROUP BY` values.
- As soon as MySQL has sent the required number of rows to the client, it aborts the query unless you are using `SQL_CALC_FOUND_ROWS`. In that case, the number of rows can be retrieved with `SELECT FOUND_ROWS()`. See Section 12.14, “Information Functions”.
- `LIMIT 0` quickly returns an empty set. This can be useful for checking the validity of a query. It can also be employed to obtain the types of the result columns within applications that use a MySQL API that makes result set metadata available. With the `mysql` client program, you can use the `--column-type-info` option to display result column types.
- If the server uses temporary tables to resolve a query, it uses the `LIMIT row_count` clause to calculate how much space is required.
- If an index is not used for `ORDER BY` but a `LIMIT` clause is also present, the optimizer may be able to avoid using a merge file and sort the rows in memory using an in-memory `filesort` operation. For details, see The In-Memory filesort Algorithm.

If multiple rows have identical values in the `ORDER BY` columns, the server is free to return those rows in any order, and may do so differently depending on the overall execution plan. In other words, the sort order of those rows is nondeterministic with respect to the nonordered columns.

One factor that affects the execution plan is `LIMIT`, so an `ORDER BY` query with and without `LIMIT` may return rows in different orders. Consider this query, which is sorted by the `category` column but nondeterministic with respect to the `id` and `rating` columns:

```
```

```
mysql> SELECT * FROM ratings ORDER BY category;
+----+----------+--------+
| id | category | rating |
+----+----------+--------+
|  1 |        1 |    4.5 |
|  5 |        1 |    3.2 |
|  3 |        2 |    3.7 |
|  4 |        2 |    3.5 |
|  6 |        2 |    3.5 |
|  2 |        3 |    5.0 |
|  7 |        3 |    2.7 |
+----+----------+--------+
```

Including `LIMIT` may affect order of rows within each `category` value. For example, this is a valid query result:

```
mysql> SELECT * FROM ratings ORDER BY category LIMIT 5;
+----+----------+--------+
| id | category | rating |
+----+----------+--------+
|  1 |        1 |    4.5 |
|  5 |        1 |    3.2 |
|  4 |        2 |    3.5 |
|  3 |        2 |    3.7 |
|  6 |        2 |    3.5 |
+----+----------+--------+
```

In each case, the rows are sorted by the `ORDER BY` column, which is all that is required by the SQL standard.

If it is important to ensure the same row order with and without `LIMIT`, include additional columns in the `ORDER BY` clause to make the order deterministic. For example, if `id` values are unique, you can make rows for a given `category` value appear in `id` order by sorting like this:

```
mysql> SELECT * FROM ratings ORDER BY category, id;
+----+----------+--------+
| id | category | rating |
+----+----------+--------+
|  1 |        1 |    4.5 |
|  5 |        1 |    3.2 |
|  3 |        2 |    3.7 |
|  4 |        2 |    3.5 |
|  6 |        2 |    3.5 |
|  2 |        3 |    5.0 |
|  7 |        3 |    2.7 |
+----+----------+--------+

mysql> SELECT * FROM ratings ORDER BY category, id LIMIT 5;
+----+----------+--------+
| id | category | rating |
+----+----------+--------+
|  1 |        1 |    4.5 |
|  5 |        1 |    3.2 |
|  3 |        2 |    3.7 |
|  4 |        2 |    3.5 |
|  6 |        2 |    3.5 |
+----+----------+--------+
```

8.2.1.17 Function Call Optimization

MySQL functions are tagged internally as deterministic or nondeterministic. A function is nondeterministic if, given fixed values for its arguments, it can return different results for different invocations. Examples of nondeterministic functions: `RAND()`, `UUID()`.

If a function is tagged nondeterministic, a reference to it in a `WHERE` clause is evaluated for every row (when selecting from one table) or combination of rows (when selecting from a multiple-table join).

MySQL also determines when to evaluate functions based on types of arguments, whether the arguments are table columns or constant values. A deterministic function that takes a table column as argument must be evaluated whenever that column changes value.

Nondeterministic functions may affect query performance. For example, some optimizations may not be available, or more locking might be required. The following discussion uses `RAND()` but applies to other nondeterministic functions as well.

Suppose that a table `t` has this definition:

```
CREATE TABLE t (id INT NOT NULL PRIMARY KEY, col_a VARCHAR(100));
```

Consider these two queries:

```
SELECT * FROM t WHERE id = POW(1,2);
SELECT * FROM t WHERE id = FLOOR(1 + RAND() * 49);
```

Both queries appear to use a primary key lookup because of the equality comparison against the primary key, but that is true only for the first of them:

- The first query always produces a maximum of one row because `POW()` with constant arguments is a constant value and is used for index lookup.
- The second query contains an expression that uses the nondeterministic function `RAND()`, which is not constant in the query but in fact has a new value for every row of table `t`. Consequently, the query reads every row of the table, evaluates the predicate for each row, and outputs all rows for which the primary key matches the random value. This might be zero, one, or multiple rows, depending on the `id` column values and the values in the `RAND()` sequence.

The effects of nondeterminism are not limited to `SELECT` statements. This `UPDATE` statement uses a nondeterministic function to select rows to be modified:

```
UPDATE t SET col_a = some_expr WHERE id = FLOOR(1 + RAND() * 49);
```

Presumably the intent is to update at most a single row for which the primary key matches the expression. However, it might update zero, one, or multiple rows, depending on the `id` column values and the values in the `RAND()` sequence.

The behavior just described has implications for performance and replication:

- Because a nondeterministic function does not produce a constant value, the optimizer cannot use strategies that might otherwise be applicable, such as index lookups. The result may be a table scan.
- `InnoDB` might escalate to a range-key lock rather than taking a single row lock for one matching row.
- Updates that do not execute deterministically are unsafe for replication.

The difficulties stem from the fact that the `RAND()` function is evaluated once for every row of the table. To avoid multiple function evaluations, use one of these techniques:

- Move the expression containing the nondeterministic function to a separate statement, saving the value in a variable. In the original statement, replace the expression with a reference to the variable, which the optimizer can treat as a constant value:

  ```
  SET @keyval = FLOOR(1 + RAND() * 49);
  UPDATE t SET col_a = some_expr WHERE id = @keyval;
  ```

- Assign the random value to a variable in a derived table. This technique causes the variable to be assigned a value, once, prior to its use in the comparison in the `WHERE` clause:

```
UPDATE /*+ NO_MERGE(dt) */ t, (SELECT FLOOR(1 + RAND() * 49) AS r) AS dt
SET col_a = some_expr WHERE id = dt.r;
```

As mentioned previously, a nondeterministic expression in the `WHERE` clause might prevent optimizations and result in a table scan. However, it may be possible to partially optimize the `WHERE` clause if other expressions are deterministic. For example:

```
SELECT * FROM t WHERE partial_key=5 AND some_column=RAND();
```

If the optimizer can use `partial_key` to reduce the set of rows selected, `RAND()` is executed fewer times, which diminishes the effect of nondeterminism on optimization.

8.2.1.18 Row Constructor Expression Optimization

Row constructors permit simultaneous comparisons of multiple values. For example, these two statements are semantically equivalent:

```
SELECT * FROM t1 WHERE (column1,column2) = (1,1);
SELECT * FROM t1 WHERE column1 = 1 AND column2 = 1;
```

In addition, the optimizer handles both expressions the same way.

The optimizer is less likely to use available indexes if the row constructor columns do not cover the prefix of an index. Consider the following table, which has a primary key on `(c1, c2, c3)`:

```
CREATE TABLE t1 (
  c1 INT, c2 INT, c3 INT, c4 CHAR(100),
  PRIMARY KEY(c1,c2,c3)
);
```

In this query, the `WHERE` clause uses all columns in the index. However, the row constructor itself does not cover an index prefix, with the result that the optimizer uses only `c1` (`key_len=4`, the size of `c1`):

```
mysql> EXPLAIN SELECT * FROM t1
       WHERE c1=1 AND (c2,c3) > (1,1)\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
   partitions: NULL
         type: ref
possible_keys: PRIMARY
          key: PRIMARY
      key_len: 4
          ref: const
         rows: 3
     filtered: 100.00
        Extra: Using where
```

In such cases, rewriting the row constructor expression using an equivalent nonconstructor expression may result in more complete index use. For the given query, the row constructor and equivalent nonconstructor expressions are:

```
(c2,c3) > (1,1)
c2 > 1 OR ((c2 = 1) AND (c3 > 1))
```

Rewriting the query to use the nonconstructor expression results in the optimizer using all three columns in the index (`key_len=12`):

```
mysql> EXPLAIN SELECT * FROM t1
```

```
      WHERE c1 = 1 AND (c2 > 1 OR ((c2 = 1) AND (c3 > 1)))\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
   partitions: NULL
         type: range
possible_keys: PRIMARY
          key: PRIMARY
      key_len: 12
          ref: NULL
         rows: 3
     filtered: 100.00
        Extra: Using where
```

Thus, for better results, avoid mixing row constructors with `AND`/`OR` expressions. Use one or the other.

Under certain conditions, the optimizer can apply the range access method to `IN()` expressions that have row constructor arguments. See Range Optimization of Row Constructor Expressions.

8.2.1.19 Avoiding Full Table Scans

The output from `EXPLAIN` shows `ALL` in the `type` column when MySQL uses a full table scan to resolve a query. This usually happens under the following conditions:

- The table is so small that it is faster to perform a table scan than to bother with a key lookup. This is common for tables with fewer than 10 rows and a short row length.
- There are no usable restrictions in the `ON` or `WHERE` clause for indexed columns.
- You are comparing indexed columns with constant values and MySQL has calculated (based on the index tree) that the constants cover too large a part of the table and that a table scan would be faster. See Section 8.2.1.1, “WHERE Clause Optimization”.
- You are using a key with low cardinality (many rows match the key value) through another column. In this case, MySQL assumes that by using the key it probably will do many key lookups and that a table scan would be faster.

For small tables, a table scan often is appropriate and the performance impact is negligible. For large tables, try the following techniques to avoid having the optimizer incorrectly choose a table scan:

- Use `ANALYZE TABLE` *`tbl_name`* to update the key distributions for the scanned table. See Section 13.7.2.1, “ANALYZE TABLE Syntax”.
- Use `FORCE INDEX` for the scanned table to tell MySQL that table scans are very expensive compared to using the given index:

  ```
  SELECT * FROM t1, t2 FORCE INDEX (index_for_column)
    WHERE t1.col_name=t2.col_name;
  ```

  See Section 8.9.4, “Index Hints”.
- Start `mysqld` with the `--max-seeks-for-key=1000` option or use `SET max_seeks_for_key=1000` to tell the optimizer to assume that no key scan causes more than 1,000 key seeks. See Section 5.1.5, “Server System Variables”.

8.2.2 Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions

The MySQL query optimizer has different strategies available to evaluate subqueries. For `IN` (or `=ANY`) subqueries, the optimizer has these choices:

- Semi-join

- Materialization
- `EXISTS` strategy

For `NOT IN` (or `<>ALL`) subqueries, the optimizer has these choices:

- Materialization
- `EXISTS` strategy

For derived tables (subqueries in the `FROM` clause), the optimizer has these choices:

- Merge the derived table into the outer query block
- Materialize the derived table to an internal temporary table

For view references and common table expressions, the optimizer has the same choices as for derived tables.

The following discussion provides more information about the preceding optimization strategies.

Note

A limitation on `UPDATE` and `DELETE` statements that use a subquery to modify a single table is that the optimizer does not use semi-join or materialization subquery optimizations. As a workaround, try rewriting them as multiple-table `UPDATE` and `DELETE` statements that use a join rather than a subquery.

8.2.2.1 Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations

The optimizer uses semi-join strategies to improve subquery execution, as described in this section.

For an inner join between two tables, the join returns a row from one table as many times as there are matches in the other table. But for some questions, the only information that matters is whether there is a match, not the number of matches. Suppose that there are tables named `class` and `roster` that list classes in a course curriculum and class rosters (students enrolled in each class), respectively. To list the classes that actually have students enrolled, you could use this join:

```
SELECT class.class_num, class.class_name
FROM class INNER JOIN roster
WHERE class.class_num = roster.class_num;
```

However, the result lists each class once for each enrolled student. For the question being asked, this is unnecessary duplication of information.

Assuming that `class_num` is a primary key in the `class` table, duplicate suppression is possible by using `SELECT DISTINCT`, but it is inefficient to generate all matching rows first only to eliminate duplicates later.

The same duplicate-free result can be obtained by using a subquery:

```
SELECT class_num, class_name
FROM class
WHERE class_num IN (SELECT class_num FROM roster);
```

Here, the optimizer can recognize that the `IN` clause requires the subquery to return only one instance of each class number from the `roster` table. In this case, the query can use a *semi-join*; that is, an operation that returns only one instance of each row in `class` that is matched by rows in `roster`.

Outer join and inner join syntax is permitted in the outer query specification, and table references may be base tables, derived tables, view references, or common table expressions.

In MySQL, a subquery must satisfy these criteria to be handled as a semi-join:

- It must be an `IN` (or `=ANY`) subquery that appears at the top level of the `WHERE` or `ON` clause, possibly as a term in an `AND` expression. For example:

  ```
  SELECT ...
  FROM ot1, ...
  WHERE (oe1, ...) IN (SELECT ie1, ... FROM it1, ... WHERE ...);
  ```

  Here, *`ot_i`* and *`it_i`* represent tables in the outer and inner parts of the query, and *`oe_i`* and *`ie_i`* represent expressions that refer to columns in the outer and inner tables.

- It must be a single `SELECT` without `UNION` constructs.
- It must not contain a `GROUP BY` or `HAVING` clause.
- It must not be implicitly grouped (it must contain no aggregate functions).
- It must not have `ORDER BY` with `LIMIT`.
- The statement must not use the `STRAIGHT_JOIN` join type in the outer query.
- The `STRAIGHT_JOIN` modifier must not be present.
- The number of outer and inner tables together must be less than the maximum number of tables permitted in a join.

The subquery may be correlated or uncorrelated. `DISTINCT` is permitted, as is `LIMIT` unless `ORDER BY` is also used.

If a subquery meets the preceding criteria, MySQL converts it to a semi-join and makes a cost-based choice from these strategies:

- Convert the subquery to a join, or use table pullout and run the query as an inner join between subquery tables and outer tables. Table pullout pulls a table out from the subquery to the outer query.
- Duplicate Weedout: Run the semi-join as if it was a join and remove duplicate records using a temporary table.
- FirstMatch: When scanning the inner tables for row combinations and there are multiple instances of a given value group, choose one rather than returning them all. This "shortcuts" scanning and eliminates production of unnecessary rows.
- LooseScan: Scan a subquery table using an index that enables a single value to be chosen from each subquery's value group.
- Materialize the subquery into an indexed temporary table that is used to perform a join, where the index is used to remove duplicates. The index might also be used later for lookups when joining the temporary table with the outer tables; if not, the table is scanned. For more information about materialization, see Section 8.2.2.2, “Optimizing Subqueries with Materialization”.

Each of these strategies can be enabled or disabled using the following `optimizer_switch` system variable flags:

- The `semijoin` flag controls whether semi-joins are used.
- If `semijoin` is enabled, the `firstmatch`, `loosescan`, `duplicateweedout`, and `materialization` flags enable finer control over the permitted semi-join strategies.
- If the `duplicateweedout` semi-join strategy is disabled, it is not used unless all other applicable strategies are also disabled.

- If `duplicateweedout` is disabled, on occasion the optimizer may generate a query plan that is far from optimal. This occurs due to heuristic pruning during greedy search, which can be avoided by setting `optimizer_prune_level=0`.

These flags are enabled by default. See Section 8.9.3, “Switchable Optimizations”.

The optimizer minimizes differences in handling of views and derived tables (subqueries in the `FROM` clause). This affects queries that use the `STRAIGHT_JOIN` modifier and a view with an `IN` subquery that can be converted to a semi-join. The following query illustrates this because the change in processing causes a change in transformation, and thus a different execution strategy:

```
CREATE VIEW v AS
SELECT *
FROM t1
WHERE a IN (SELECT b
            FROM t2);

SELECT STRAIGHT_JOIN *
FROM t3 JOIN v ON t3.x = v.a;
```

The optimizer first looks at the view and converts the `IN` subquery to a semi-join, then checks whether it is possible to merge the view into the outer query. Because the `STRAIGHT_JOIN` modifier in the outer query prevents semi-join, the optimizer refuses the merge, causing derived table evaluation using a materialized table.

`EXPLAIN` output indicates the use of semi-join strategies as follows:

- For extended `EXPLAIN` output, the text displayed by a following `SHOW WARNINGS` shows the rewritten query, which displays the semi-join structure. (See Section 8.8.3, “Extended EXPLAIN Output Format”.) From this you can get an idea about which tables were pulled out of the semi-join. If a subquery was converted to a semi-join, you will see that the subquery predicate is gone and its tables and `WHERE` clause were merged into the outer query join list and `WHERE` clause.
- Temporary table use for Duplicate Weedout is indicated by `Start temporary` and `End temporary` in the `Extra` column. Tables that were not pulled out and are in the range of `EXPLAIN` output rows covered by `Start temporary` and `End temporary` have their `rowid` in the temporary table.
- `FirstMatch(`*`tbl_name`*`)` in the `Extra` column indicates join shortcutting.
- `LooseScan(`*`m`*`..`*`n`*`)` in the `Extra` column indicates use of the LooseScan strategy. *`m`* and *`n`* are key part numbers.
- Temporary table use for materialization is indicated by rows with a `select_type` value of `MATERIALIZED` and rows with a `table` value of `<subquery`*`N`*`>`.

8.2.2.2 Optimizing Subqueries with Materialization

The optimizer uses materialization to enable more efficient subquery processing. Materialization speeds up query execution by generating a subquery result as a temporary table, normally in memory. The first time MySQL needs the subquery result, it materializes that result into a temporary table. Any subsequent time the result is needed, MySQL refers again to the temporary table. The optimizer may index the table with a hash index to make lookups fast and inexpensive. The index is unique, which eliminates duplicates and makes the table smaller.

Subquery materialization uses an in-memory temporary table when possible, falling back to on-disk storage if the table becomes too large. See Section 8.4.4, “Internal Temporary Table Use in MySQL”.

If materialization is not used, the optimizer sometimes rewrites a noncorrelated subquery as a correlated subquery. For example, the following `IN` subquery is noncorrelated (*`where_condition`* involves only columns from `t2` and not `t1`):

```
SELECT * FROM t1
WHERE t1.a IN (SELECT t2.b FROM t2 WHERE where_condition);
```

The optimizer might rewrite this as an EXISTS correlated subquery:

```
SELECT * FROM t1
WHERE EXISTS (SELECT t2.b FROM t2 WHERE where_condition AND t1.a=t2.b);
```

Subquery materialization using a temporary table avoids such rewrites and makes it possible to execute the subquery only once rather than once per row of the outer query.

For subquery materialization to be used in MySQL, the optimizer_switch system variable materialization flag must be enabled. (See Section 8.9.3, “Switchable Optimizations”.) With the materialization flag enabled, materialization applies to subquery predicates that appear anywhere (in the select list, WHERE, ON, GROUP BY, HAVING, or ORDER BY), for predicates that fall into any of these use cases:

- The predicate has this form, when no outer expression *oe_i* or inner expression *ie_i* is nullable. *N* is 1 or larger.

  ```
  (oe_1, oe_2, ..., oe_N) [NOT] IN (SELECT ie_1, i_2, ..., ie_N ...)
  ```

- The predicate has this form, when there is a single outer expression *oe* and inner expression *ie*. The expressions can be nullable.

  ```
  oe [NOT] IN (SELECT ie ...)
  ```

- The predicate is IN or NOT IN and a result of UNKNOWN (NULL) has the same meaning as a result of FALSE.

The following examples illustrate how the requirement for equivalence of UNKNOWN and FALSE predicate evaluation affects whether subquery materialization can be used. Assume that *where_condition* involves columns only from t2 and not t1 so that the subquery is noncorrelated.

This query is subject to materialization:

```
SELECT * FROM t1
WHERE t1.a IN (SELECT t2.b FROM t2 WHERE where_condition);
```

Here, it does not matter whether the IN predicate returns UNKNOWN or FALSE. Either way, the row from t1 is not included in the query result.

An example where subquery materialization is not used is the following query, where t2.b is a nullable column:

```
SELECT * FROM t1
WHERE (t1.a,t1.b) NOT IN (SELECT t2.a,t2.b FROM t2
                          WHERE where_condition);
```

The following restrictions apply to the use of subquery materialization:

- The types of the inner and outer expressions must match. For example, the optimizer might be able to use materialization if both expressions are integer or both are decimal, but cannot if one expression is integer and the other is decimal.
- The inner expression cannot be a BLOB.

Use of EXPLAIN with a query provides some indication of whether the optimizer uses subquery materialization. Compared to query execution that does not use materialization, select_type may change from DEPENDENT SUBQUERY to SUBQUERY. This indicates that, for a subquery that would

be executed once per outer row, materialization enables the subquery to be executed just once. In addition, for extended `EXPLAIN` output, the text displayed by a following `SHOW WARNINGS` includes `materialize` and `materialized-subquery`.

8.2.2.3 Optimizing Derived Tables, View References, and Common Table Expressions

The optimizer can handle derived tables (subqueries in the `FROM` clause) references using two strategies:

- Merge the derived table into the outer query block
- Materialize the derived table to an internal temporary table

The optimizer uses the same strategies to handle view references and common table expressions.

Example 1:

```
SELECT * FROM (SELECT * FROM t1) AS derived_t1;
```

With merging, that query is executed similar to:

```
SELECT * FROM t1;
```

Example 2:

```
SELECT *
  FROM t1 JOIN (SELECT t2.f1 FROM t2) AS derived_t2 ON t1.f2=derived_t2.f1
  WHERE t1.f1 > 0;
```

With merging, that query is executed similar to:

```
SELECT t1.*, t2.f1
  FROM t1 JOIN t2 ON t1.f2=t2.f1
  WHERE t1.f1 > 0;
```

With materialization, `derived_t1` and `derived_t2` are treated as a separate table within their respective queries.

The optimizer handles derived tables and view references the same way: It avoids unnecessary materialization whenever possible, which enables pushing down conditions from the outer query to derived tables and produces more efficient execution plans. (For an example, see Section 8.2.2.2, “Optimizing Subqueries with Materialization”.)

If merging would result in an outer query block that references more than 61 base tables, the optimizer chooses materialization instead.

The optimizer propagates an `ORDER BY` clause in a derived table or view reference to the outer query block if these conditions are all true:

- The outer query is not grouped or aggregated.
- The outer query does not specify `DISTINCT`, `HAVING`, or `ORDER BY`.
- The outer query has this derived table or view reference as the only source in the `FROM` clause.

Otherwise, the optimizer ignores the `ORDER BY` clause.

The following means are available to influence whether the optimizer attempts to merge derived tables and view references into the outer query block:

- The `MERGE` and `NO_MERGE` optimizer hints can be used. They apply assuming that no other rule prevents merging. See Section 8.9.2, “Optimizer Hints”.

- Similarly, you can use the `derived_merge` flag of the `optimizer_switch` system variable. See Section 8.9.3, “Switchable Optimizations”. By default, the flag is enabled to permit merging. Disabling the flag prevents merging and avoids `ER_UPDATE_TABLE_USED` errors.

  The `derived_merge` flag also applies to views that contain no `ALGORITHM` clause. Thus, if an `ER_UPDATE_TABLE_USED` error occurs for a view reference that uses an expression equivalent to the subquery, adding `ALGORITHM=TEMPTABLE` to the view definition prevents merging and takes precedence over the `derived_merge` value.

- It is possible to disable merging by using in the subquery any constructs that prevent merging, although these are not as explicit in their effect on materialization. Constructs that prevent merging are the same for derived tables and view references:

  - Aggregate functions (`SUM()`, `MIN()`, `MAX()`, `COUNT()`, and so forth)
  - `DISTINCT`
  - `GROUP BY`
  - `HAVING`
  - `LIMIT`
  - `UNION` or `UNION ALL`
  - Subqueries in the select list
  - Assignments to user variables
  - Refererences only to literal values (in this case, there is no underlying table)

If the optimizer chooses the materialization strategy rather than merging for a derived table, it handles the query as follows:

- The optimizer postpones derived table materialization until its contents are needed during query execution. This improves performance because delaying materialization may result in not having to do it at all. Consider a query that joins the result of a derived table to another table: If the optimizer processes that other table first and finds that it returns no rows, the join need not be carried out further and the optimizer can completely skip materializing the derived table.
- During query execution, the optimizer may add an index to a derived table to speed up row retrieval from it.

Consider the following `EXPLAIN` statement, for which a subquery appears in the `FROM` clause of a `SELECT` query:

```
EXPLAIN SELECT * FROM (SELECT * FROM t1) AS derived_t1;
```

The optimizer avoids materializing the subquery by delaying it until the result is needed during `SELECT` execution. In this case, the query is not executed (because it occurs in an `EXPLAIN` statement), so the result is never needed.

Even for queries that are executed, delay of subquery materialization may enable the optimizer to avoid materialization entirely. When this happens, query execution is quicker by the time needed to perform materialization. Consider the following query, which joins the result of a subquery in the `FROM` clause to another table:

```
SELECT *
  FROM t1 JOIN (SELECT t2.f1 FROM t2) AS derived_t2
          ON t1.f2=derived_t2.f1
  WHERE t1.f1 > 0;
```

If the optimization processes `t1` first and the `WHERE` clause produces an empty result, the join must necessarily be empty and the subquery need not be materialized.

For cases when a derived table requires materialization, the optimizer may add an index to the materialized table to speed up access to it. If such an index enables `ref` access to the table, it can greatly reduce amount of data read during query execution. Consider the following query:

```
SELECT *
 FROM t1 JOIN (SELECT DISTINCT f1 FROM t2) AS derived_t2
         ON t1.f1=derived_t2.f1;
```

The optimizer constructs an index over column `f1` from `derived_t2` if doing so would enable use of `ref` access for the lowest cost execution plan. After adding the index, the optimizer can treat the materialized derived table the same as a regular table with an index, and it benefits similarly from the generated index. The overhead of index creation is negligible compared to the cost of query execution without the index. If `ref` access would result in higher cost than some other access method, the optimizer creates no index and loses nothing.

For optimizer trace output, a merged derived table or view reference is not shown as a node. Only its underlying tables appear in the top query's plan.

What is true for materialization of derived tables is also true for common table expressions (CTEs). In addition, the following considerations pertain specifically to CTEs.

If a CTE is materialized by a query, it is materialized once for the query, even if the query references it several times.

A recursive CTE is always materialized.

If a CTE is materialized, the optimizer automatically adds relevant indexes if it estimates that indexing will speed up access by the top-level statement to the CTE. This is similar to automatic indexing of derived tables, except that if the CTE is referenced multiple times, the optimizer may create multiple indexes, to speed up access by each reference in the most appropriate way.

The `MERGE` and `NO_MERGE` optimizer hints can be applied to CTEs. Each CTE reference in the top-level statement can have its own hint, permitting CTE references to be selectively merged or materialized. The following statement uses hints to indicate that `cte1` should be merged and `cte2` should be materialized:

```
WITH
  cte1 AS (SELECT a, b FROM table1),
  cte2 AS (SELECT c, d FROM table2)
SELECT /*+ MERGE(cte1) NO_MERGE(cte2) */ cte1.b, cte2.d
FROM cte1 JOIN cte2
WHERE cte1.a = cte2.c;
```

The `ALGORITHM` clause for `CREATE VIEW` does not affect materialization for any `WITH` clause preceding the `SELECT` statement in the view definition. Consider this statement:

```
CREATE ALGORITHM={TEMPTABLE|MERGE} VIEW v1 AS WITH ... SELECT ...
```

The `ALGORITHM` value affects materialization only of the `SELECT`, not the `WITH` clause.

For CTEs, the storage engine used for on-disk internal temporary tables cannot be `MyISAM`. If `internal_tmp_disk_storage_engine=MYISAM`, an error occurs for any attempt to materialize a CTE using an on-disk temporary table.

As mentioned previously, a CTE, if materialized, is materialized once, even if referenced multiple times. To indicate one-time materialization, optimizer trace output contains an occurrence of `creating_tmp_table` plus one or more occurrences of `reusing_tmp_table`.

CTEs are similar to derived tables, for which the `materialized_from_subquery` node follows the reference. This is true for a CTE that is referenced multiple times, so there is no duplication of `materialized_from_subquery` nodes (which would give the impression that the subquery is executed multiple times, and produce unnecessarily verbose output). Only one reference to the CTE has a complete `materialized_from_subquery` node with the description of its subquery plan. Other references have a reduced `materialized_from_subquery` node. The same idea applies to `EXPLAIN` output in `TRADITIONAL` format: Subqueries for other references are not shown.

8.2.2.4 Optimizing Subqueries with the EXISTS Strategy

Certain optimizations are applicable to comparisons that use the `IN` (or `=ANY`) operator to test subquery results. This section discusses these optimizations, particularly with regard to the challenges that `NULL` values present. The last part of the discussion suggests how you can help the optimizer.

Consider the following subquery comparison:

```
outer_expr IN (SELECT inner_expr FROM ... WHERE subquery_where)
```

MySQL evaluates queries “from outside to inside.” That is, it first obtains the value of the outer expression *`outer_expr`*, and then runs the subquery and captures the rows that it produces.

A very useful optimization is to “inform” the subquery that the only rows of interest are those where the inner expression *`inner_expr`* is equal to *`outer_expr`*. This is done by pushing down an appropriate equality into the subquery's `WHERE` clause to make it more restrictive. The converted comparison looks like this:

```
EXISTS (SELECT 1 FROM ... WHERE subquery_where AND outer_expr=inner_expr)
```

After the conversion, MySQL can use the pushed-down equality to limit the number of rows it must examine to evaluate the subquery.

More generally, a comparison of *`N`* values to a subquery that returns *`N`*-value rows is subject to the same conversion. If *`oe_i`* and *`ie_i`* represent corresponding outer and inner expression values, this subquery comparison:

```
(oe_1, ..., oe_N) IN
  (SELECT ie_1, ..., ie_N FROM ... WHERE subquery_where)
```

Becomes:

```
EXISTS (SELECT 1 FROM ... WHERE subquery_where
                          AND oe_1 = ie_1
                          AND ...
                          AND oe_N = ie_N)
```

For simplicity, the following discussion assumes a single pair of outer and inner expression values.

The conversion just described has its limitations. It is valid only if we ignore possible `NULL` values. That is, the “pushdown” strategy works as long as both of these conditions are true:

- *`outer_expr`* and *`inner_expr`* cannot be `NULL`.
- You need not distinguish `NULL` from `FALSE` subquery results. If the subquery is a part of an `OR` or `AND` expression in the `WHERE` clause, MySQL assumes that you do not care. Another instance where the optimizer notices that `NULL` and `FALSE` subquery results need not be distinguished is this construct:

  ```
  ... WHERE outer_expr IN (subquery)
  ```

In this case, the `WHERE` clause rejects the row whether `IN (subquery)` returns `NULL` or `FALSE`.

When either or both of those conditions do not hold, optimization is more complex.

Suppose that *`outer_expr`* is known to be a non-`NULL` value but the subquery does not produce a row such that *`outer_expr`* = *`inner_expr`*. Then *`outer_expr`* `IN (SELECT ...)` evaluates as follows:

- `NULL`, if the `SELECT` produces any row where *`inner_expr`* is `NULL`
- `FALSE`, if the `SELECT` produces only non-`NULL` values or produces nothing

In this situation, the approach of looking for rows with *`outer_expr`* = *`inner_expr`* is no longer valid. It is necessary to look for such rows, but if none are found, also look for rows where *`inner_expr`* is `NULL`. Roughly speaking, the subquery can be converted to something like this:

```
EXISTS (SELECT 1 FROM ... WHERE subquery_where AND
        (outer_expr=inner_expr OR inner_expr IS NULL))
```

The need to evaluate the extra `IS NULL` condition is why MySQL has the `ref_or_null` access method:

```
mysql> EXPLAIN
       SELECT outer_expr IN (SELECT t2.maybe_null_key
                             FROM t2, t3 WHERE ...)
       FROM t1;
*************************** 1. row ***************************
           id: 1
  select_type: PRIMARY
        table: t1
...
*************************** 2. row ***************************
           id: 2
  select_type: DEPENDENT SUBQUERY
        table: t2
         type: ref_or_null
possible_keys: maybe_null_key
          key: maybe_null_key
      key_len: 5
          ref: func
         rows: 2
        Extra: Using where; Using index
...
```

The `unique_subquery` and `index_subquery` subquery-specific access methods also have “or `NULL`” variants.

The additional `OR ... IS NULL` condition makes query execution slightly more complicated (and some optimizations within the subquery become inapplicable), but generally this is tolerable.

The situation is much worse when *`outer_expr`* can be `NULL`. According to the SQL interpretation of `NULL` as “unknown value,” `NULL IN (SELECT` *`inner_expr`* `...)` should evaluate to:

- `NULL`, if the `SELECT` produces any rows
- `FALSE`, if the `SELECT` produces no rows

For proper evaluation, it is necessary to be able to check whether the `SELECT` has produced any rows at all, so *`outer_expr`* = *`inner_expr`* cannot be pushed down into the subquery. This is a problem because many real world subqueries become very slow unless the equality can be pushed down.

Essentially, there must be different ways to execute the subquery depending on the value of *`outer_expr`*.

The optimizer chooses SQL compliance over speed, so it accounts for the possibility that *outer_expr* might be NULL:

- If *outer_expr* is NULL, to evaluate the following expression, it is necessary to execute the SELECT to determine whether it produces any rows:

  ```
  NULL IN (SELECT inner_expr FROM ... WHERE subquery_where)
  ```

  It is necessary to execute the original SELECT here, without any pushed-down equalities of the kind mentioned earlier.

- On the other hand, when *outer_expr* is not NULL, it is absolutely essential that this comparison:

  ```
  outer_expr IN (SELECT inner_expr FROM ... WHERE subquery_where)
  ```

  Be converted to this expression that uses a pushed-down condition:

  ```
  EXISTS (SELECT 1 FROM ... WHERE subquery_where AND outer_expr=inner_expr)
  ```

  Without this conversion, subqueries will be slow.

To solve the dilemma of whether or not to push down conditions into the subquery, the conditions are wrapped within “trigger” functions. Thus, an expression of the following form:

```
outer_expr IN (SELECT inner_expr FROM ... WHERE subquery_where)
```

Is converted into:

```
EXISTS (SELECT 1 FROM ... WHERE subquery_where
                          AND trigcond(outer_expr=inner_expr))
```

More generally, if the subquery comparison is based on several pairs of outer and inner expressions, the conversion takes this comparison:

```
(oe_1, ..., oe_N) IN (SELECT ie_1, ..., ie_N FROM ... WHERE subquery_where)
```

And converts it to this expression:

```
EXISTS (SELECT 1 FROM ... WHERE subquery_where
                          AND trigcond(oe_1=ie_1)
                          AND ...
                          AND trigcond(oe_N=ie_N)
       )
```

Each `trigcond(X)` is a special function that evaluates to the following values:

- *X* when the “linked” outer expression *oe_i* is not NULL
- TRUE when the “linked” outer expression *oe_i* is NULL

Note

Trigger functions are *not* triggers of the kind that you create with CREATE TRIGGER.

Equalities that are wrapped within `trigcond()` functions are not first class predicates for the query optimizer. Most optimizations cannot deal with predicates that may be turned on and off at query execution time, so they assume any `trigcond(X)` to be an unknown function and ignore it. Triggered equalities can be used by those optimizations:

- Reference optimizations: `trigcond(X=Y [OR Y IS NULL])` can be used to construct `ref`, `eq_ref`, or `ref_or_null` table accesses.
- Index lookup-based subquery execution engines: `trigcond(X=Y)` can be used to construct `unique_subquery` or `index_subquery` accesses.
- Table-condition generator: If the subquery is a join of several tables, the triggered condition is checked as soon as possible.

When the optimizer uses a triggered condition to create some kind of index lookup-based access (as for the first two items of the preceding list), it must have a fallback strategy for the case when the condition is turned off. This fallback strategy is always the same: Do a full table scan. In `EXPLAIN` output, the fallback shows up as `Full scan on NULL key` in the `Extra` column:

```
mysql> EXPLAIN SELECT t1.col1,
       t1.col1 IN (SELECT t2.key1 FROM t2 WHERE t2.col2=t1.col2) FROM t1\G
*************************** 1. row ***************************
           id: 1
  select_type: PRIMARY
        table: t1
        ...
*************************** 2. row ***************************
           id: 2
  select_type: DEPENDENT SUBQUERY
        table: t2
         type: index_subquery
possible_keys: key1
          key: key1
      key_len: 5
          ref: func
         rows: 2
        Extra: Using where; Full scan on NULL key
```

If you run `EXPLAIN` followed by `SHOW WARNINGS`, you can see the triggered condition:

```
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: select `test`.`t1`.`col1` AS `col1`,
         <in_optimizer>(`test`.`t1`.`col1`,
         <exists>(<index_lookup>(<cache>(`test`.`t1`.`col1`) in t2
         on key1 checking NULL
         where (`test`.`t2`.`col2` = `test`.`t1`.`col2`) having
         trigcond(<is_not_null_test>(`test`.`t2`.`key1`))))) AS
         `t1.col1 IN (select t2.key1 from t2 where t2.col2=t1.col2)`
         from `test`.`t1`
```

The use of triggered conditions has some performance implications. A `NULL IN (SELECT ...)` expression now may cause a full table scan (which is slow) when it previously did not. This is the price paid for correct results (the goal of the trigger-condition strategy is to improve compliance, not speed).

For multiple-table subqueries, execution of `NULL IN (SELECT ...)` is particularly slow because the join optimizer does not optimize for the case where the outer expression is `NULL`. It assumes that subquery evaluations with `NULL` on the left side are very rare, even if there are statistics that indicate otherwise. On the other hand, if the outer expression might be `NULL` but never actually is, there is no performance penalty.

To help the query optimizer better execute your queries, use these suggestions:

- Declare a column as `NOT NULL` if it really is. This also helps other aspects of the optimizer by simplifying condition testing for the column.
- If you need not distinguish a `NULL` from `FALSE` subquery result, you can easily avoid the slow execution path. Replace a comparison that looks like this:

```
outer_expr IN (SELECT inner_expr FROM ...)
```

with this expression:

```
(outer_expr IS NOT NULL) AND (outer_expr IN (SELECT inner_expr FROM ...))
```

Then `NULL IN (SELECT ...)` is never evaluated because MySQL stops evaluating `AND` parts as soon as the expression result is clear.

Another possible rewrite:

```
EXISTS (SELECT inner_expr FROM ...
        WHERE inner_expr=outer_expr)
```

This would apply when you need not distinguish `NULL` from `FALSE` subquery results, in which case you may actually want `EXISTS`.

The `subquery_materialization_cost_based` flag of the `optimizer_switch` system variable enables control over the choice between subquery materialization and `IN`-to-`EXISTS` subquery transformation. See Section 8.9.3, “Switchable Optimizations”.

8.2.3 Optimizing INFORMATION_SCHEMA Queries

Applications that monitor databases may make frequent use of `INFORMATION_SCHEMA` tables. To write queries for these tables most efficiently, use the following general guidelines:

- Try to query only `INFORMATION_SCHEMA` tables that are views on data dictionary tables.
- Try to query only for static metadata. Selecting columns or using retrieval conditions for dynamic metadata along with static metadata adds overhead to process the dynamic metadata.

Note

Comparison behavior for database and table names in `INFORMATION_SCHEMA` queries might differ from what you expect. For details, see Section 10.1.8.7, “Using Collation in INFORMATION_SCHEMA Searches”.

These `INFORMATION_SCHEMA` tables are implemented as views on data dictionary tables, so queries on them retrieve information from the data dictionary:

```
CHARACTER_SETS
COLLATIONS
COLLATION_CHARACTER_SET_APPLICABILITY
COLUMNS
EVENTS
KEY_COLUMN_USAGE
PARAMETERS
ROUTINES
SCHEMATA
STATISTICS
TABLES
TABLE_CONSTRAINTS
TRIGGERS
VIEWS
```

Some types of values, even for a non-view `INFORMATION_SCHEMA` table, are retrieved by lookups from the data dictionary. This includes values such as database and table names, table types, and storage engines.

Some `INFORMATION_SCHEMA` tables contain columns that provide table statistics:

```
STATISTICS.CARDINALITY
TABLES.AUTO_INCREMENT
TABLES.AVG_ROW_LENGTH
TABLES.CHECKSUM
TABLES.CHECK_TIME
TABLES.CREATE_TIME
TABLES.DATA_FREE
TABLES.DATA_LENGTH
TABLES.INDEX_LENGTH
TABLES.MAX_DATA_LENGTH
TABLES.TABLE_ROWS
TABLES.UPDATE_TIME
```

Those columns represent dynamic table metadata; that is, information that changes as table contents change.

The server has two sources from which to retrieve table statistics. The `information_schema_stats` system variable controls which table statistics source the server uses:

- When `information_schema_stats` is `CACHED` (the default), the server uses cached statistics stored in the `STATISTICS` and `TABLES` tables.
- When `information_schema_stats` is `LATEST`, the server obtains statistics directly from storage engines.

To determine whether to set `information_schema_stats` to `CACHED` or `LATEST`, consider the tradeoffs:

- When the server starts, cached statistics are `NULL`. To update them for a given table, use `ANALYZE TABLE`. This incurs a one-time statistics-calculation cost, but the cached statistics remain up to date to the extent that the table changes slowly. To update the cached statistics at any time thereafter, use `ANALYZE TABLE` again.
- For use of the latest statistics, `ANALYZE TABLE` is not needed, but each query that retrieves statistics incurs somewhat higher execution cost than for use of cached statistics.

Set the global `information_schema_stats` value to determine the default used initially by all sessions. Individual sessions can set the session `information_schema_stats` value to override the global value as desired.

Note

If the `innodb_read_only` system variable is enabled, `ANALYZE TABLE` may fail because it cannot update statistics tables in the data dictionary, which use `InnoDB`. For `ANALYZE TABLE` operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a `MyISAM` table). To obtain the updated distribution statistics, set `information_schema_stats=LATEST`.

For `INFORMATION_SCHEMA` tables implemented as views on data dictionary tables, indexes on the underlying data dictionary tables permit the optimizer to construct efficient query execution plans. To see the choices made by the optimizer, use `EXPLAIN`. To also see the query used by the server to execute an `INFORMATION_SCHEMA` query, use `SHOW WARNINGS` immediately following `EXPLAIN`.

Consider this statement, which identifies collations for the `utf8mb4` character set:

```
mysql> SELECT COLLATION_NAME
       FROM INFORMATION_SCHEMA.COLLATION_CHARACTER_SET_APPLICABILITY
       WHERE CHARACTER_SET_NAME = 'utf8mb4';
+----------------------------+
| COLLATION_NAME             |
+----------------------------+
| utf8mb4_general_ci         |
```

```
| utf8mb4_bin                |
| utf8mb4_unicode_ci         |
| utf8mb4_icelandic_ci       |
| utf8mb4_latvian_ci         |
| utf8mb4_romanian_ci        |
| utf8mb4_slovenian_ci       |
...
```

How does the server process that statement? To find out, use `EXPLAIN`:

```
mysql> EXPLAIN SELECT COLLATION_NAME
       FROM INFORMATION_SCHEMA.COLLATION_CHARACTER_SET_APPLICABILITY
       WHERE CHARACTER_SET_NAME = 'utf8mb4'\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: cs
   partitions: NULL
         type: const
possible_keys: PRIMARY,name
          key: name
      key_len: 194
          ref: const
         rows: 1
     filtered: 100.00
        Extra: Using index
*************************** 2. row ***************************
           id: 1
  select_type: SIMPLE
        table: col
   partitions: NULL
         type: ref
possible_keys: character_set_id
          key: character_set_id
      key_len: 8
          ref: const
         rows: 68
     filtered: 100.00
        Extra: NULL
2 rows in set, 1 warning (0.01 sec)
```

To see the query used to statisfy that statement, use `SHOW WARNINGS`:

```
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: /* select#1 */ select `mysql`.`col`.`name` AS `COLLATION_NAME`
         from `mysql`.`character_sets` `cs`
         join `mysql`.`collations` `col`
         where ((`mysql`.`col`.`character_set_id` = '45')
         and ('utf8mb4' = 'utf8mb4'))
```

As indicated by `SHOW WARNINGS`, the server handles the query on `COLLATION_CHARACTER_SET_APPLICABILITY` as a query on the `character_sets` and `collations` data dictionary tables in the `mysql` system database.

8.2.4 Optimizing Performance Schema Queries

Applications that monitor databases may make frequent use of Performance Schema tables. To write queries for these tables most efficiently, take advantage of their indexes. For example, include a `WHERE` clause that restricts retrieved rows based on comparison to specific values in an indexed column.

Most Performance Schema tables have indexes. Tables that do not are those that normally contain few rows or are unlikely to be queried frequently. Performance Schema indexes give the optimizer access to execution plans other than full table scans. These indexes also improve performance for related objects, such as `sys` schema views that use those tables.

To see whether a given Performance Schema table has indexes and what they are, use `SHOW INDEX` or `SHOW CREATE TABLE`:

```
mysql> SHOW INDEX FROM performance_schema.accounts\G
*************************** 1. row ***************************
        Table: accounts
   Non_unique: 0
     Key_name: ACCOUNT
 Seq_in_index: 1
  Column_name: USER
    Collation: NULL
  Cardinality: NULL
     Sub_part: NULL
       Packed: NULL
         Null: YES
   Index_type: HASH
      Comment:
Index_comment:
      Visible: YES
*************************** 2. row ***************************
        Table: accounts
   Non_unique: 0
     Key_name: ACCOUNT
 Seq_in_index: 2
  Column_name: HOST
    Collation: NULL
  Cardinality: NULL
     Sub_part: NULL
       Packed: NULL
         Null: YES
   Index_type: HASH
      Comment:
Index_comment:
      Visible: YES

mysql> SHOW CREATE TABLE performance_schema.rwlock_instances\G
*************************** 1. row ***************************
       Table: rwlock_instances
Create Table: CREATE TABLE `rwlock_instances` (
  `NAME` varchar(128) NOT NULL,
  `OBJECT_INSTANCE_BEGIN` bigint(20) unsigned NOT NULL,
  `WRITE_LOCKED_BY_THREAD_ID` bigint(20) unsigned DEFAULT NULL,
  `READ_LOCKED_BY_COUNT` int(10) unsigned NOT NULL,
  PRIMARY KEY (`OBJECT_INSTANCE_BEGIN`),
  KEY `NAME` (`NAME`),
  KEY `WRITE_LOCKED_BY_THREAD_ID` (`WRITE_LOCKED_BY_THREAD_ID`)
) ENGINE=PERFORMANCE_SCHEMA DEFAULT CHARSET=utf8
```

To see the execution plan for a Performance Schema query and whether it uses any indexes, use `EXPLAIN`:

```
mysql> EXPLAIN SELECT * FROM performance_schema.accounts
       WHERE (USER,HOST) = ('root','localhost')\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: accounts
   partitions: NULL
         type: const
possible_keys: ACCOUNT
          key: ACCOUNT
      key_len: 278
          ref: const,const
         rows: 1
     filtered: 100.00
        Extra: NULL
```

The `EXPLAIN` output indicates that the optimizer uses the `accounts` table `ACCOUNT` index that comprises the `USER` and `HOST` columns.

Performance Schema indexes are virtual: They are a construct of the Performance Schema storage engine and use no memory or disk storage. The Performance Schema reports index information to the optimizer so that it can construct efficient execution plans. The Performance Schema in turn uses optimizer information about what to look for (for example, a particular key value), so that it can perform efficient lookups without building actual index structures. This implementation provides two important benefits:

- It entirely avoids the maintenance cost normally incurred for tables that undergo frequent updates.
- It reduces at an early stage of query execution the amount of data retrieved. For conditions on the indexed columns, the Performance Schema efficiently returns only table rows that satisfy the query conditions. Without an index, the Performance Schema would return all rows in the table, requiring that the optimizer later evaluate the conditions against each row to produce the final result.

Performance Schema indexes are predefined and cannot be dropped, added, or altered.

Performance Schema indexes are similar to hash indexes. For example:

- They are used only for equality comparisons that use the `=` or `<=>` operators.
- They are unordered. If a query result must have specific row ordering characteristics, include an `ORDER BY` clause.

For additional information about hash indexes, see Section 8.3.8, “Comparison of B-Tree and Hash Indexes”.

8.2.5 Optimizing Data Change Statements

This section explains how to speed up data change statements: `INSERT`, `UPDATE`, and `DELETE`. Traditional OLTP applications and modern web applications typically do many small data change operations, where concurrency is vital. Data analysis and reporting applications typically run data change operations that affect many rows at once, where the main considerations is the I/O to write large amounts of data and keep indexes up-to-date. For inserting and updating large volumes of data (known in the industry as ETL, for “extract-transform-load”), sometimes you use other SQL statements or external commands, that mimic the effects of `INSERT`, `UPDATE`, and `DELETE` statements.

8.2.5.1 Optimizing INSERT Statements

To optimize insert speed, combine many small operations into a single large operation. Ideally, you make a single connection, send the data for many new rows at once, and delay all index updates and consistency checking until the very end.

The time required for inserting a row is determined by the following factors, where the numbers indicate approximate proportions:

- Connecting: (3)
- Sending query to server: (2)
- Parsing query: (2)
- Inserting row: (1 × size of row)
- Inserting indexes: (1 × number of indexes)
- Closing: (1)

This does not take into consideration the initial overhead to open tables, which is done once for each concurrently running query.

The size of the table slows down the insertion of indexes by log *N*, assuming B-tree indexes.

You can use the following methods to speed up inserts:

- If you are inserting many rows from the same client at the same time, use `INSERT` statements with multiple `VALUES` lists to insert several rows at a time. This is considerably faster (many times faster in some cases) than using separate single-row `INSERT` statements. If you are adding data to a nonempty table, you can tune the `bulk_insert_buffer_size` variable to make data insertion even faster. See Section 5.1.5, “Server System Variables”.
- When loading a table from a text file, use `LOAD DATA INFILE`. This is usually 20 times faster than using `INSERT` statements. See Section 13.2.7, “LOAD DATA INFILE Syntax”.
- Take advantage of the fact that columns have default values. Insert values explicitly only when the value to be inserted differs from the default. This reduces the parsing that MySQL must do and improves the insert speed.
- See Section 8.5.5, “Bulk Data Loading for InnoDB Tables” for tips specific to `InnoDB` tables.
- See Section 8.6.2, “Bulk Data Loading for MyISAM Tables” for tips specific to `MyISAM` tables.

8.2.5.2 Optimizing UPDATE Statements

An update statement is optimized like a `SELECT` query with the additional overhead of a write. The speed of the write depends on the amount of data being updated and the number of indexes that are updated. Indexes that are not changed do not get updated.

Another way to get fast updates is to delay updates and then do many updates in a row later. Performing multiple updates together is much quicker than doing one at a time if you lock the table.

For a `MyISAM` table that uses dynamic row format, updating a row to a longer total length may split the row. If you do this often, it is very important to use `OPTIMIZE TABLE` occasionally. See Section 13.7.2.4, “OPTIMIZE TABLE Syntax”.

8.2.5.3 Optimizing DELETE Statements

The time required to delete individual rows in a `MyISAM` table is exactly proportional to the number of indexes. To delete rows more quickly, you can increase the size of the key cache by increasing the `key_buffer_size` system variable. See Section 5.1.1, “Configuring the Server”.

To delete all rows from a `MyISAM` table, `TRUNCATE TABLE tbl_name` is faster than `DELETE FROM tbl_name`. Truncate operations are not transaction-safe; an error occurs when attempting one in the course of an active transaction or active table lock. See Section 13.1.30, “TRUNCATE TABLE Syntax”.

8.2.6 Optimizing Database Privileges

The more complex your privilege setup, the more overhead applies to all SQL statements. Simplifying the privileges established by `GRANT` statements enables MySQL to reduce permission-checking overhead when clients execute statements. For example, if you do not grant any table-level or column-level privileges, the server need not ever check the contents of the `tables_priv` and `columns_priv` tables. Similarly, if you place no resource limits on any accounts, the server does not have to perform resource counting. If you have a very high statement-processing load, consider using a simplified grant structure to reduce permission-checking overhead.

8.2.7 Other Optimization Tips

This section lists a number of miscellaneous tips for improving query processing speed:

- If your application makes several database requests to perform related updates, combining the statements into a stored routine can help performance. Similarly, if your application computes a single result based on several column values or large volumes of data, combining the computation into a UDF (user-defined function) can help performance. The resulting fast database operations are then available to be reused by other queries, applications, and even code written in different

programming languages. See Section 23.2, “Using Stored Routines (Procedures and Functions)” and Section 28.4, “Adding New Functions to MySQL” for more information.

- To fix any compression issues that occur with `ARCHIVE` tables, use `OPTIMIZE TABLE`. See Section 16.5, “The ARCHIVE Storage Engine”.
- If possible, classify reports as “live” or as “statistical”, where data needed for statistical reports is created only from summary tables that are generated periodically from the live data.
- If you have data that does not conform well to a rows-and-columns table structure, you can pack and store data into a `BLOB` column. In this case, you must provide code in your application to pack and unpack information, but this might save I/O operations to read and write the sets of related values.
- With Web servers, store images and other binary assets as files, with the path name stored in the database rather than the file itself. Most Web servers are better at caching files than database contents, so using files is generally faster. (Although you must handle backups and storage issues yourself in this case.)
- If you need really high speed, look at the low-level MySQL interfaces. For example, by accessing the MySQL `InnoDB` or `MyISAM` storage engine directly, you could get a substantial speed increase compared to using the SQL interface.
- Replication can provide a performance benefit for some operations. You can distribute client retrievals among replication servers to split up the load. To avoid slowing down the master while making backups, you can make backups using a slave server. See Chapter 18, *Replication*.

8.3 Optimization and Indexes

The best way to improve the performance of `SELECT` operations is to create indexes on one or more of the columns that are tested in the query. The index entries act like pointers to the table rows, allowing the query to quickly determine which rows match a condition in the `WHERE` clause, and retrieve the other column values for those rows. All MySQL data types can be indexed.

Although it can be tempting to create an indexes for every possible column used in a query, unnecessary indexes waste space and waste time for MySQL to determine which indexes to use. Indexes also add to the cost of inserts, updates, and deletes because each index must be updated. You must find the right balance to achieve fast queries using the optimal set of indexes.

8.3.1 How MySQL Uses Indexes

Indexes are used to find rows with specific column values quickly. Without an index, MySQL must begin with the first row and then read through the entire table to find the relevant rows. The larger the table, the more this costs. If the table has an index for the columns in question, MySQL can quickly determine the position to seek to in the middle of the data file without having to look at all the data. This is much faster than reading every row sequentially.

Most MySQL indexes (`PRIMARY KEY`, `UNIQUE`, `INDEX`, and `FULLTEXT`) are stored in B-trees. Exceptions: Indexes on spatial data types use R-trees; `MEMORY` tables also support hash indexes; `InnoDB` uses inverted lists for `FULLTEXT` indexes.

In general, indexes are used as described in the following discussion. Characteristics specific to hash indexes (as used in `MEMORY` tables) are described in Section 8.3.8, “Comparison of B-Tree and Hash Indexes”.

MySQL uses indexes for these operations:

- To find the rows matching a `WHERE` clause quickly.
- To eliminate rows from consideration. If there is a choice between multiple indexes, MySQL normally uses the index that finds the smallest number of rows (the most selective index).

- If the table has a multiple-column index, any leftmost prefix of the index can be used by the optimizer to look up rows. For example, if you have a three-column index on `(col1, col2, col3)`, you have indexed search capabilities on `(col1)`, `(col1, col2)`, and `(col1, col2, col3)`. For more information, see Section 8.3.5, “Multiple-Column Indexes”.

- To retrieve rows from other tables when performing joins. MySQL can use indexes on columns more efficiently if they are declared as the same type and size. In this context, `VARCHAR` and `CHAR` are considered the same if they are declared as the same size. For example, `VARCHAR(10)` and `CHAR(10)` are the same size, but `VARCHAR(10)` and `CHAR(15)` are not.

  For comparisons between nonbinary string columns, both columns should use the same character set. For example, comparing a `utf8` column with a `latin1` column precludes use of an index.

  Comparison of dissimilar columns (comparing a string column to a temporal or numeric column, for example) may prevent use of indexes if values cannot be compared directly without conversion. For a given value such as `1` in the numeric column, it might compare equal to any number of values in the string column such as `'1'`, `' 1'`, `'00001'`, or `'01.e1'`. This rules out use of any indexes for the string column.

- To find the `MIN()` or `MAX()` value for a specific indexed column *`key_col`*. This is optimized by a preprocessor that checks whether you are using `WHERE key_part_N = constant` on all key parts that occur before *`key_col`* in the index. In this case, MySQL does a single key lookup for each `MIN()` or `MAX()` expression and replaces it with a constant. If all expressions are replaced with constants, the query returns at once. For example:

```
SELECT MIN(key_part2),MAX(key_part2)
  FROM tbl_name WHERE key_part1=10;
```

- To sort or group a table if the sorting or grouping is done on a leftmost prefix of a usable index (for example, `ORDER BY key_part1, key_part2`). If all key parts are followed by `DESC`, the key is read in reverse order. (Or, if the index is a descending index, the key is read in forward order.) See Section 8.2.1.13, “ORDER BY Optimization”, Section 8.2.1.14, “GROUP BY Optimization”, and Section 8.3.12, “Descending Indexes”.

- In some cases, a query can be optimized to retrieve values without consulting the data rows. (An index that provides all the necessary results for a query is called a covering index.) If a query uses from a table only columns that are included in some index, the selected values can be retrieved from the index tree for greater speed:

```
SELECT key_part3 FROM tbl_name
  WHERE key_part1=1
```

Indexes are less important for queries on small tables, or big tables where report queries process most or all of the rows. When a query needs to access most of the rows, reading sequentially is faster than working through an index. Sequential reads minimize disk seeks, even if not all the rows are needed for the query. See Section 8.2.1.19, “Avoiding Full Table Scans” for details.

8.3.2 Using Primary Keys

The primary key for a table represents the column or set of columns that you use in your most vital queries. It has an associated index, for fast query performance. Query performance benefits from the `NOT NULL` optimization, because it cannot include any `NULL` values. With the `InnoDB` storage engine, the table data is physically organized to do ultra-fast lookups and sorts based on the primary key column or columns.

If your table is big and important, but does not have an obvious column or set of columns to use as a primary key, you might create a separate column with auto-increment values to use as the primary key. These unique IDs can serve as pointers to corresponding rows in other tables when you join tables using foreign keys.

8.3.3 Using Foreign Keys

If a table has many columns, and you query many different combinations of columns, it might be efficient to split the less-frequently used data into separate tables with a few columns each, and relate them back to the main table by duplicating the numeric ID column from the main table. That way, each small table can have a primary key for fast lookups of its data, and you can query just the set of columns that you need using a join operation. Depending on how the data is distributed, the queries might perform less I/O and take up less cache memory because the relevant columns are packed together on disk. (To maximize performance, queries try to read as few data blocks as possible from disk; tables with only a few columns can fit more rows in each data block.)

8.3.4 Column Indexes

The most common type of index involves a single column, storing copies of the values from that column in a data structure, allowing fast lookups for the rows with the corresponding column values. The B-tree data structure lets the index quickly find a specific value, a set of values, or a range of values, corresponding to operators such as `=`, `>`, `≤`, `BETWEEN`, `IN`, and so on, in a `WHERE` clause.

The maximum number of indexes per table and the maximum index length is defined per storage engine. See Chapter 15, *The InnoDB Storage Engine*, and Chapter 16, *Alternative Storage Engines*. All storage engines support at least 16 indexes per table and a total index length of at least 256 bytes. Most storage engines have higher limits.

For additional information about column indexes, see Section 13.1.12, “CREATE INDEX Syntax”.

- Index Prefixes
- FULLTEXT Indexes
- Spatial Indexes
- Indexes in the MEMORY Storage Engine

Index Prefixes

With *`col_name(N)`* syntax in an index specification for a string column, you can create an index that uses only the first *`N`* characters of the column. Indexing only a prefix of column values in this way can make the index file much smaller. When you index a `BLOB` or `TEXT` column, you *must* specify a prefix length for the index. For example:

```
CREATE TABLE test (blob_col BLOB, INDEX(blob_col(10)));
```

Prefixes can be up to 767 bytes long for `InnoDB` tables that use the `REDUNDANT` or `COMPACT` row format. The length limit is raised to 3072 bytes for `InnoDB` tables that use the `DYNAMIC` or `COMPRESSED` row format. For MyISAM tables, the prefix limit is 1000 bytes.

Note

Prefix limits are measured in bytes, whereas the prefix length in `CREATE TABLE`, `ALTER TABLE`, and `CREATE INDEX` statements is interpreted as number of characters for nonbinary string types (`CHAR`, `VARCHAR`, `TEXT`) and number of bytes for binary string types (`BINARY`, `VARBINARY`, `BLOB`). Take this into account when specifying a prefix length for a nonbinary string column that uses a multibyte character set.

For additional information about index prefixes, see Section 13.1.12, “CREATE INDEX Syntax”.

FULLTEXT Indexes

`FULLTEXT` indexes are used for full-text searches. Only the `InnoDB` and `MyISAM` storage engines support `FULLTEXT` indexes and only for `CHAR`, `VARCHAR`, and `TEXT` columns. Indexing always takes

place over the entire column and column prefix indexing is not supported. For details, see Section 12.9, “Full-Text Search Functions”.

Optimizations are applied to certain kinds of `FULLTEXT` queries against single `InnoDB` tables. Queries with these characteristics are particularly efficient:

- `FULLTEXT` queries that only return the document ID, or the document ID and the search rank.
- `FULLTEXT` queries that sort the matching rows in descending order of score and apply a `LIMIT` clause to take the top N matching rows. For this optimization to apply, there must be no `WHERE` clauses and only a single `ORDER BY` clause in descending order.
- `FULLTEXT` queries that retrieve only the `COUNT(*)` value of rows matching a search term, with no additional `WHERE` clauses. Code the `WHERE` clause as `WHERE MATCH(text) AGAINST ('other_text')`, without any `> 0` comparison operator.

For queries that contain full-text expressions, MySQL evaluates those expressions during the optimization phase of query execution. The optimizer does not just look at full-text expressions and make estimates, it actually evaluates them in the process of developing an execution plan.

An implication of this behavior is that `EXPLAIN` for full-text queries is typically slower than for non-full-text queries for which no expression evaluation occurs during the optimization phase.

`EXPLAIN` for full-text queries may show `Select tables optimized away` in the `Extra` column due to matching occurring during optimization; in this case, no table access need occur during later execution.

Spatial Indexes

You can create indexes on spatial data types. `MyISAM` and `InnoDB` support R-tree indexes on spatial types. Other storage engines use B-trees for indexing spatial types (except for `ARCHIVE`, which does not support spatial type indexing).

Indexes in the MEMORY Storage Engine

The `MEMORY` storage engine uses `HASH` indexes by default, but also supports `BTREE` indexes.

8.3.5 Multiple-Column Indexes

MySQL can create composite indexes (that is, indexes on multiple columns). An index may consist of up to 16 columns. For certain data types, you can index a prefix of the column (see Section 8.3.4, “Column Indexes”).

MySQL can use multiple-column indexes for queries that test all the columns in the index, or queries that test just the first column, the first two columns, the first three columns, and so on. If you specify the columns in the right order in the index definition, a single composite index can speed up several kinds of queries on the same table.

A multiple-column index can be considered a sorted array, the rows of which contain values that are created by concatenating the values of the indexed columns.

Note

As an alternative to a composite index, you can introduce a column that is “hashed” based on information from other columns. If this column is short, reasonably unique, and indexed, it might be faster than a “wide” index on many columns. In MySQL, it is very easy to use this extra column:

```
SELECT * FROM tbl_name
  WHERE hash_col=MD5(CONCAT(val1,val2))
  AND col1=val1 AND col2=val2;
```

Suppose that a table has the following specification:

```
CREATE TABLE test (
    id         INT NOT NULL,
    last_name  CHAR(30) NOT NULL,
    first_name CHAR(30) NOT NULL,
    PRIMARY KEY (id),
    INDEX name (last_name,first_name)
);
```

The `name` index is an index over the `last_name` and `first_name` columns. The index can be used for lookups in queries that specify values in a known range for combinations of `last_name` and `first_name` values. It can also be used for queries that specify just a `last_name` value because that column is a leftmost prefix of the index (as described later in this section). Therefore, the `name` index is used for lookups in the following queries:

```
SELECT * FROM test WHERE last_name='Widenius';

SELECT * FROM test
  WHERE last_name='Widenius' AND first_name='Michael';

SELECT * FROM test
  WHERE last_name='Widenius'
  AND (first_name='Michael' OR first_name='Monty');

SELECT * FROM test
  WHERE last_name='Widenius'
  AND first_name >='M' AND first_name < 'N';
```

However, the `name` index is *not* used for lookups in the following queries:

```
SELECT * FROM test WHERE first_name='Michael';

SELECT * FROM test
  WHERE last_name='Widenius' OR first_name='Michael';
```

Suppose that you issue the following `SELECT` statement:

```
SELECT * FROM tbl_name
  WHERE col1=val1 AND col2=val2;
```

If a multiple-column index exists on `col1` and `col2`, the appropriate rows can be fetched directly. If separate single-column indexes exist on `col1` and `col2`, the optimizer attempts to use the Index Merge optimization (see Section 8.2.1.3, “Index Merge Optimization”), or attempts to find the most restrictive index by deciding which index excludes more rows and using that index to fetch the rows.

If the table has a multiple-column index, any leftmost prefix of the index can be used by the optimizer to look up rows. For example, if you have a three-column index on `(col1, col2, col3)`, you have indexed search capabilities on `(col1)`, `(col1, col2)`, and `(col1, col2, col3)`.

MySQL cannot use the index to perform lookups if the columns do not form a leftmost prefix of the index. Suppose that you have the `SELECT` statements shown here:

```
SELECT * FROM tbl_name WHERE col1=val1;
SELECT * FROM tbl_name WHERE col1=val1 AND col2=val2;

SELECT * FROM tbl_name WHERE col2=val2;
SELECT * FROM tbl_name WHERE col2=val2 AND col3=val3;
```

If an index exists on `(col1, col2, col3)`, only the first two queries use the index. The third and fourth queries do involve indexed columns, but `(col2)` and `(col2, col3)` are not leftmost prefixes of `(col1, col2, col3)`.

8.3.6 Verifying Index Usage

Always check whether all your queries really use the indexes that you have created in the tables. Use the `EXPLAIN` statement, as described in Section 8.8.1, “Optimizing Queries with EXPLAIN”.

8.3.7 InnoDB and MyISAM Index Statistics Collection

Storage engines collect statistics about tables for use by the optimizer. Table statistics are based on value groups, where a value group is a set of rows with the same key prefix value. For optimizer purposes, an important statistic is the average value group size.

MySQL uses the average value group size in the following ways:

- To estimate how may rows must be read for each `ref` access
- To estimate how many row a partial join will produce; that is, the number of rows that an operation of this form will produce:

```
(...) JOIN tbl_name ON tbl_name.key = expr
```

As the average value group size for an index increases, the index is less useful for those two purposes because the average number of rows per lookup increases: For the index to be good for optimization purposes, it is best that each index value target a small number of rows in the table. When a given index value yields a large number of rows, the index is less useful and MySQL is less likely to use it.

The average value group size is related to table cardinality, which is the number of value groups. The `SHOW INDEX` statement displays a cardinality value based on *`N/S`*, where *`N`* is the number of rows in the table and *`S`* is the average value group size. That ratio yields an approximate number of value groups in the table.

For a join based on the `<=>` comparison operator, `NULL` is not treated differently from any other value: `NULL <=> NULL`, just as *`N <=> N`* for any other *`N`*.

However, for a join based on the `=` operator, `NULL` is different from non-`NULL` values: *`expr1 = expr2`* is not true when *`expr1`* or *`expr2`* (or both) are `NULL`. This affects `ref` accesses for comparisons of the form *`tbl_name.key = expr`*: MySQL will not access the table if the current value of *`expr`* is `NULL`, because the comparison cannot be true.

For `=` comparisons, it does not matter how many `NULL` values are in the table. For optimization purposes, the relevant value is the average size of the non-`NULL` value groups. However, MySQL does not currently enable that average size to be collected or used.

For `InnoDB` and `MyISAM` tables, you have some control over collection of table statistics by means of the `innodb_stats_method` and `myisam_stats_method` system variables, respectively. These variables have three possible values, which differ as follows:

- When the variable is set to `nulls_equal`, all `NULL` values are treated as identical (that is, they all form a single value group).

  If the `NULL` value group size is much higher than the average non-`NULL` value group size, this method skews the average value group size upward. This makes index appear to the optimizer to be less useful than it really is for joins that look for non-`NULL` values. Consequently, the `nulls_equal` method may cause the optimizer not to use the index for `ref` accesses when it should.

- When the variable is set to `nulls_unequal`, `NULL` values are not considered the same. Instead, each `NULL` value forms a separate value group of size 1.

  If you have many `NULL` values, this method skews the average value group size downward. If the average non-`NULL` value group size is large, counting `NULL` values each as a group of size 1 causes the optimizer to overestimate the value of the index for joins that look for non-`NULL` values.

Consequently, the `nulls_unequal` method may cause the optimizer to use this index for `ref` lookups when other methods may be better.

- When the variable is set to `nulls_ignored`, `NULL` values are ignored.

If you tend to use many joins that use `<=>` rather than `=`, `NULL` values are not special in comparisons and one `NULL` is equal to another. In this case, `nulls_equal` is the appropriate statistics method.

The `innodb_stats_method` system variable has a global value; the `myisam_stats_method` system variable has both global and session values. Setting the global value affects statistics collection for tables from the corresponding storage engine. Setting the session value affects statistics collection only for the current client connection. This means that you can force a table's statistics to be regenerated with a given method without affecting other clients by setting the session value of `myisam_stats_method`.

To regenerate `MyISAM` table statistics, you can use any of the following methods:

- Execute `myisamchk --stats_method=method_name --analyze`
- Change the table to cause its statistics to go out of date (for example, insert a row and then delete it), and then set `myisam_stats_method` and issue an `ANALYZE TABLE` statement

Some caveats regarding the use of `innodb_stats_method` and `myisam_stats_method`:

- You can force table statistics to be collected explicitly, as just described. However, MySQL may also collect statistics automatically. For example, if during the course of executing statements for a table, some of those statements modify the table, MySQL may collect statistics. (This may occur for bulk inserts or deletes, or some `ALTER TABLE` statements, for example.) If this happens, the statistics are collected using whatever value `innodb_stats_method` or `myisam_stats_method` has at the time. Thus, if you collect statistics using one method, but the system variable is set to the other method when a table's statistics are collected automatically later, the other method will be used.
- There is no way to tell which method was used to generate statistics for a given table.
- These variables apply only to `InnoDB` and `MyISAM` tables. Other storage engines have only one method for collecting table statistics. Usually it is closer to the `nulls_equal` method.

8.3.8 Comparison of B-Tree and Hash Indexes

Understanding the B-tree and hash data structures can help predict how different queries perform on different storage engines that use these data structures in their indexes, particularly for the `MEMORY` storage engine that lets you choose B-tree or hash indexes.

- B-Tree Index Characteristics
- Hash Index Characteristics

B-Tree Index Characteristics

A B-tree index can be used for column comparisons in expressions that use the `=`, `>`, `>=`, `<`, `<=`, or `BETWEEN` operators. The index also can be used for `LIKE` comparisons if the argument to `LIKE` is a constant string that does not start with a wildcard character. For example, the following `SELECT` statements use indexes:

```
SELECT * FROM tbl_name WHERE key_col LIKE 'Patrick%';
SELECT * FROM tbl_name WHERE key_col LIKE 'Pat%_ck%';
```

In the first statement, only rows with `'Patrick' <= key_col < 'Patricl'` are considered. In the second statement, only rows with `'Pat' <= key_col < 'Pau'` are considered.

The following `SELECT` statements do not use indexes:

```
SELECT * FROM tbl_name WHERE key_col LIKE '%Patrick%';
SELECT * FROM tbl_name WHERE key_col LIKE other_col;
```

In the first statement, the `LIKE` value begins with a wildcard character. In the second statement, the `LIKE` value is not a constant.

If you use `... LIKE '%string%'` and *`string`* is longer than three characters, MySQL uses the *Turbo Boyer-Moore algorithm* to initialize the pattern for the string and then uses this pattern to perform the search more quickly.

A search using *`col_name`* `IS NULL` employs indexes if *`col_name`* is indexed.

Any index that does not span all `AND` levels in the `WHERE` clause is not used to optimize the query. In other words, to be able to use an index, a prefix of the index must be used in every `AND` group.

The following `WHERE` clauses use indexes:

```
... WHERE index_part1=1 AND index_part2=2 AND other_column=3

    /* index = 1 OR index = 2 */
... WHERE index=1 OR A=10 AND index=2

    /* optimized like "index_part1='hello'" */
... WHERE index_part1='hello' AND index_part3=5

    /* Can use index on index1 but not on index2 or index3 */
... WHERE index1=1 AND index2=2 OR index1=3 AND index3=3;
```

These `WHERE` clauses do *not* use indexes:

```
    /* index_part1 is not used */
... WHERE index_part2=1 AND index_part3=2

    /*  Index is not used in both parts of the WHERE clause  */
... WHERE index=1 OR A=10

    /* No index spans all rows  */
... WHERE index_part1=1 OR index_part2=10
```

Sometimes MySQL does not use an index, even if one is available. One circumstance under which this occurs is when the optimizer estimates that using the index would require MySQL to access a very large percentage of the rows in the table. (In this case, a table scan is likely to be much faster because it requires fewer seeks.) However, if such a query uses `LIMIT` to retrieve only some of the rows, MySQL uses an index anyway, because it can much more quickly find the few rows to return in the result.

Hash Index Characteristics

Hash indexes have somewhat different characteristics from those just discussed:

- They are used only for equality comparisons that use the `=` or `<=>` operators (but are *very* fast). They are not used for comparison operators such as `<` that find a range of values. Systems that rely on this type of single-value lookup are known as “key-value stores”; to use MySQL for such applications, use hash indexes wherever possible.
- The optimizer cannot use a hash index to speed up `ORDER BY` operations. (This type of index cannot be used to search for the next entry in order.)
- MySQL cannot determine approximately how many rows there are between two values (this is used by the range optimizer to decide which index to use). This may affect some queries if you change a `MyISAM` or `InnoDB` table to a hash-indexed `MEMORY` table.
- Only whole keys can be used to search for a row. (With a B-tree index, any leftmost prefix of the key can be used to find rows.)

8.3.9 Use of Index Extensions

InnoDB automatically extends each secondary index by appending the primary key columns to it. Consider this table definition:

```
CREATE TABLE t1 (
  i1 INT NOT NULL DEFAULT 0,
  i2 INT NOT NULL DEFAULT 0,
  d DATE DEFAULT NULL,
  PRIMARY KEY (i1, i2),
  INDEX k_d (d)
) ENGINE = InnoDB;
```

This table defines the primary key on columns (i1, i2). It also defines a secondary index k_d on column (d), but internally InnoDB extends this index and treats it as columns (d, i1, i2).

The optimizer takes into account the primary key columns of the extended secondary index when determining how and whether to use that index. This can result in more efficient query execution plans and better performance.

The optimizer can use extended secondary indexes for ref, range, and index_merge index access, for loose index scans, for join and sorting optimization, and for MIN()/MAX() optimization.

The following example shows how execution plans are affected by whether the optimizer uses extended secondary indexes. Suppose that t1 is populated with these rows:

```
INSERT INTO t1 VALUES
(1, 1, '1998-01-01'), (1, 2, '1999-01-01'),
(1, 3, '2000-01-01'), (1, 4, '2001-01-01'),
(1, 5, '2002-01-01'), (2, 1, '1998-01-01'),
(2, 2, '1999-01-01'), (2, 3, '2000-01-01'),
(2, 4, '2001-01-01'), (2, 5, '2002-01-01'),
(3, 1, '1998-01-01'), (3, 2, '1999-01-01'),
(3, 3, '2000-01-01'), (3, 4, '2001-01-01'),
(3, 5, '2002-01-01'), (4, 1, '1998-01-01'),
(4, 2, '1999-01-01'), (4, 3, '2000-01-01'),
(4, 4, '2001-01-01'), (4, 5, '2002-01-01'),
(5, 1, '1998-01-01'), (5, 2, '1999-01-01'),
(5, 3, '2000-01-01'), (5, 4, '2001-01-01'),
(5, 5, '2002-01-01');
```

Now consider this query:

```
EXPLAIN SELECT COUNT(*) FROM t1 WHERE i1 = 3 AND d = '2000-01-01'
```

The optimizer cannot use the primary key in this case because that comprises columns (i1, i2) and the query does not refer to i2. Instead, the optimizer can use the secondary index k_d on (d), and the execution plan depends on whether the extended index is used.

When the optimizer does not consider index extensions, it treats the index k_d as only (d). EXPLAIN for the query produces this result:

```
mysql> EXPLAIN SELECT COUNT(*) FROM t1 WHERE i1 = 3 AND d = '2000-01-01'\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
         type: ref
possible_keys: PRIMARY,k_d
          key: k_d
      key_len: 4
          ref: const
         rows: 5
        Extra: Using where; Using index
```

When the optimizer takes index extensions into account, it treats `k_d` as `(d, i1, i2)`. In this case, it can use the leftmost index prefix `(d, i1)` to produce a better execution plan:

```
mysql> EXPLAIN SELECT COUNT(*) FROM t1 WHERE i1 = 3 AND d = '2000-01-01'\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
         type: ref
possible_keys: PRIMARY,k_d
          key: k_d
      key_len: 8
          ref: const,const
         rows: 1
        Extra: Using index
```

In both cases, `key` indicates that the optimizer will use secondary index `k_d` but the `EXPLAIN` output shows these improvements from using the extended index:

- `key_len` goes from 4 bytes to 8 bytes, indicating that key lookups use columns `d` and `i1`, not just `d`.
- The `ref` value changes from `const` to `const,const` because the key lookup uses two key parts, not one.
- The `rows` count decreases from 5 to 1, indicating that `InnoDB` should need to examine fewer rows to produce the result.
- The `Extra` value changes from `Using where; Using index` to `Using index`. This means that rows can be read using only the index, without consulting columns in the data row.

Differences in optimizer behavior for use of extended indexes can also be seen with `SHOW STATUS`:

```
FLUSH TABLE t1;
FLUSH STATUS;
SELECT COUNT(*) FROM t1 WHERE i1 = 3 AND d = '2000-01-01';
SHOW STATUS LIKE 'handler_read%'
```

The preceding statements include `FLUSH TABLE` and `FLUSH STATUS` to flush the table cache and clear the status counters.

Without index extensions, `SHOW STATUS` produces this result:

```
+-----------------------+-------+
| Variable_name         | Value |
+-----------------------+-------+
| Handler_read_first    | 0     |
| Handler_read_key      | 1     |
| Handler_read_last     | 0     |
| Handler_read_next     | 5     |
| Handler_read_prev     | 0     |
| Handler_read_rnd      | 0     |
| Handler_read_rnd_next | 0     |
+-----------------------+-------+
```

With index extensions, `SHOW STATUS` produces this result. The `Handler_read_next` value decreases from 5 to 1, indicating more efficient use of the index:

```
+-----------------------+-------+
| Variable_name         | Value |
+-----------------------+-------+
| Handler_read_first    | 0     |
| Handler_read_key      | 1     |
| Handler_read_last     | 0     |
| Handler_read_next     | 1     |
| Handler_read_prev     | 0     |
```

```
| Handler_read_rnd      | 0     |
| Handler_read_rnd_next | 0     |
+-----------------------+-------+
```

The `use_index_extensions` flag of the `optimizer_switch` system variable permits control over whether the optimizer takes the primary key columns into account when determining how to use an `InnoDB` table's secondary indexes. By default, `use_index_extensions` is enabled. To check whether disabling use of index extensions will improve performance, use this statement:

```
SET optimizer_switch = 'use_index_extensions=off';
```

Use of index extensions by the optimizer is subject to the usual limits on the number of key parts in an index (16) and the maximum key length (3072 bytes).

8.3.10 Optimizer Use of Generated Column Indexes

MySQL supports indexes on generated columns. For example:

```
CREATE TABLE t1 (f1 INT, gc INT AS (f1 + 1) STORED, INDEX (gc));
```

The generated column, `gc`, is defined as the expression `f1 + 1`. The column is also indexed and the optimizer can take that index into account during execution plan construction. In the following query, the `WHERE` clause refers to `gc` and the optimizer considers whether the index on that column yields a more efficient plan:

```
SELECT * FROM t1 WHERE gc > 9;
```

The optimizer can use indexes on generated columns to generate execution plans, even in the absence of direct references in queries to those columns by name. This occurs if the `WHERE`, `ORDER BY`, or `GROUP BY` clause refers to an expression that matches the definition of some indexed generated column. The following query does not refer directly to `gc` but does use an expression that matches the definition of `gc`:

```
SELECT * FROM t1 WHERE f1 + 1 > 9;
```

The optimizer recognizes that the expression `f1 + 1` matches the definition of `gc` and that `gc` is indexed, so it considers that index during execution plan construction. You can see this using `EXPLAIN`:

```
mysql> EXPLAIN SELECT * FROM t1 WHERE f1 + 1 > 9\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: t1
   partitions: NULL
         type: range
possible_keys: gc
          key: gc
      key_len: 5
          ref: NULL
         rows: 1
     filtered: 100.00
        Extra: Using index condition
```

In effect, the optimizer has replaced the expression `f1 + 1` with the name of the generated column that matches the expression. That is also apparent in the rewritten query available in the extended `EXPLAIN` information displayed by `SHOW WARNINGS`:

```
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
```

```
   Code: 1003
Message: /* select#1 */ select `test`.`t1`.`f1` AS `f1`,`test`.`t1`.`gc`
         AS `gc` from `test`.`t1` where (`test`.`t1`.`gc` > 9)
```

The following restrictions and conditions apply to the optimizer's use of generated column indexes:

- For a query expression to match a generated column definition, the expression must be identical and it must have the same result type. For example, if the generated column expression is `f1 + 1`, the optimizer will not recognize a match if the query uses `1 + f1`, or if `f1 + 1` (an integer expression) is compared with a string.
- The optimization applies to these operators: `=`, `<`, `<=`, `>`, `>=`, `BETWEEN`, and `IN()`.

  For operators other than `BETWEEN` and `IN()`, either operand can be replaced by a matching generated column. For `BETWEEN` and `IN()`, only the first argument can be replaced by a matching generated column, and the other arguments must have the same result type. `BETWEEN` and `IN()` are not yet supported for comparisons involving JSON values.
- The generated column must be defined as an expression that contains at least a function call or one of the operators mentioned in the preceding item. The expression cannot consist of a simple reference to another column. For example, `gc INT AS (f1) STORED` consists only of a column reference, so indexes on `gc` are not considered.
- For comparisons of strings to indexed generated columns that compute a value from a JSON function that returns a quoted string, `JSON_UNQUOTE()` is needed in the column definition to remove the extra quotes from the function value. (For direct comparison of a string to the function result, the JSON comparator handles quote removal, but this does not occur for index lookups.) For example, instead of writing a column definition like this:

  ```
  doc_name TEXT AS (JSON_EXTRACT(jdoc, '$.name')) STORED
  ```

  Write it like this:

  ```
  doc_name TEXT AS (JSON_UNQUOTE(JSON_EXTRACT(jdoc, '$.name'))) STORED
  ```

  With the latter definition, the optimizer can detect a match for both of these comparisons:

  ```
  ... WHERE JSON_EXTRACT(jdoc, '$.name') = 'some_string' ...
  ... WHERE JSON_UNQUOTE(JSON_EXTRACT(jdoc, '$.name')) = 'some_string' ...
  ```

  Without `JSON_UNQUOTE()` in the column definition, the optimizer detects a match only for the first of those comparisons.
- If the optimizer picks the wrong index, an index hint can be used to disable it and force the optimizer to make a different choice.

8.3.11 Invisible Indexes

MySQL supports invisible indexes; that is, indexes that are not used by the optimizer. The feature applies to indexes other than primary keys (either explicit or implicit). This feature is storage engine neutral (supported for any engine).

Indexes are visible by default. To control index visibility explicitly for a new index, use a `VISIBLE` or `INVISIBLE` keyword as part of the index definition for `CREATE TABLE`, `CREATE INDEX`, or `ALTER TABLE`:

```
CREATE TABLE t1 (
  i INT,
  j INT,
  k INT,
```

```
  INDEX i_idx (i) INVISIBLE
) ENGINE = InnoDB;
CREATE INDEX j_idx ON t1 (j) INVISIBLE;
ALTER TABLE t1 ADD INDEX k_idx (k) INVISIBLE;
```

To alter the visibility of an existing index, use a `VISIBLE` or `INVISIBLE` keyword with the `ALTER TABLE ... ALTER INDEX` operation:

```
ALTER TABLE t1 ALTER INDEX i_idx INVISIBLE;
ALTER TABLE t1 ALTER INDEX i_idx VISIBLE;
```

Information about whether an index is visible or invisible is available from the `INFORMATION_SCHEMA.STATISTICS` table or `SHOW INDEX` output. For example:

```
mysql> SELECT INDEX_NAME, IS_VISIBLE
       FROM INFORMATION_SCHEMA.STATISTICS
       WHERE TABLE_SCHEMA = 'db1' AND TABLE_NAME = 't1';
+------------+------------+
| INDEX_NAME | IS_VISIBLE |
+------------+------------+
| i_idx      | YES        |
| j_idx      | NO         |
| k_idx      | NO         |
+------------+------------+
```

Invisible indexes make it possible to test the effect of removing an index on query performance, without making a destructive change that must be undone should the index turn out to be required. Dropping and re-adding an index can be expensive for a large table, whereas making it invisible and visible are fast, in-place operations.

If an index made invisible actually is needed or used by the optimizer, there are several ways to notice the effect of its absence on queries for the table:

- Errors occur for queries that include index hints that refer to the invisible index.
- Performance Schema data shows an increase in workload for affected queries.
- Queries have different `EXPLAIN` execution plans.
- Queries appear in the slow query log that did not appear there previously.

Index visibility does not affect index maintenance. For example, an index continues to be updated per changes to table rows, and a unique index prevents insertion of duplicates into a column, regardless of whether the index is visible or invisible.

A table with no explicit primary key may still have an effective implicit primary key if it has any `UNIQUE` indexes on `NOT NULL` columns. In this case, the first such index places the same constraint on table rows as an explicit primary key and that index cannot be made invisible. Consider the following table definition:

```
CREATE TABLE t2 (
  i INT NOT NULL,
  j INT NOT NULL,
  UNIQUE j_idx (j)
) ENGINE = InnoDB;
```

The definition includes no explicit primary key, but the index on `NOT NULL` column `j` places the same constraint on rows as a primary key and cannot be made invisible:

```
mysql> ALTER TABLE t2 ALTER INDEX j_idx INVISIBLE;
ERROR 3522 (HY000): A primary key index cannot be invisible.
```

Now suppose that an explicit primary key is added to the table:

```
ALTER TABLE t2 ADD PRIMARY KEY (i);
```

The explicit primary key cannot be made invisible. In addition, the unique index on `j` no longer acts as an implicit primary key and as a result can be made invisible:

```
mysql> ALTER TABLE t2 ALTER INDEX j_idx INVISIBLE;
Query OK, 0 rows affected (0.03 sec)
```

8.3.12 Descending Indexes

MySQL supports descending indexes: `DESC` in an index definition is no longer ignored but causes storage of key values in descending order. Previously, indexes could be scanned in reverse order but at a performance penalty. A descending index can be scanned in forward order, which is more efficient. Descending indexes also make it possible for the optimizer to use multiple-column indexes when the most efficient scan order mixes ascending order for some columns and descending order for others.

Consider the following table definition, which contains two columns and four two-column index definitions for the various combinations of ascending and descending indexes on the columns:

```
CREATE TABLE t (
  c1 INT, c2 INT,
  INDEX idx1 (c1 ASC, c2 ASC),
  INDEX idx2 (c1 ASC, c2 DESC),
  INDEX idx3 (c1 DESC, c2 ASC),
  INDEX idx4 (c1 DESC, c2 DESC)
);
```

The table definition results in four distinct indexes. The optimizer can perform a forward index scan for each of the `ORDER BY` clauses and need not use a `filesort` operation:

```
ORDER BY c1 ASC, c2 ASC    -- optimizer can use idx1
ORDER BY c1 DESC, c2 DESC  -- optimizer can use idx4
ORDER BY c1 ASC, c2 DESC   -- optimizer can use idx2
ORDER BY c1 DESC, c2 ASC   -- optimizer can use idx3
```

Use of descending indexes is subject to these conditions:

- Descending indexes are supported only for the `InnoDB` storage engine, with these limitations:
  - Change buffering is not supported for a secondary index if the index contains a descending index key column or if the primary key includes a descending index column.
  - The `InnoDB` SQL parser does not use descending indexes. For `InnoDB` full-text search, this means that the index required on the `FTS_DOC_ID` column of the indexed table cannot be defined as a descending index. For more information, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”.
- Descending indexes are supported for all data types for which ascending indexes are available.
- Descending indexes are supported for ordinary (nongenerated) and generated columns (both `VIRTUAL` and `STORED`).
- `DISTINCT` can use any index containing matching columns, including descending key parts.
- Indexes that have descending key parts are not used for `MIN()`/`MAX()` optimization of queries that invoke aggregate functions but do not have a `GROUP BY` clause.
- Descending indexes are supported for `BTREE` but not `HASH` indexes. Descending indexes are not supported for `FULLTEXT` or `SPATIAL` indexes.

  Explicitly specified `ASC` and `DESC` designators for `HASH`, `FULLTEXT`, and `SPATIAL` indexes results in an error.

8.4 Optimizing Database Structure

In your role as a database designer, look for the most efficient way to organize your schemas, tables, and columns. As when tuning application code, you minimize I/O, keep related items together, and plan ahead so that performance stays high as the data volume increases. Starting with an efficient database design makes it easier for team members to write high-performing application code, and makes the database likely to endure as applications evolve and are rewritten.

8.4.1 Optimizing Data Size

Design your tables to minimize their space on the disk. This can result in huge improvements by reducing the amount of data written to and read from disk. Smaller tables normally require less main memory while their contents are being actively processed during query execution. Any space reduction for table data also results in smaller indexes that can be processed faster.

MySQL supports many different storage engines (table types) and row formats. For each table, you can decide which storage and indexing method to use. Choosing the proper table format for your application can give you a big performance gain. See Chapter 15, *The InnoDB Storage Engine*, and Chapter 16, *Alternative Storage Engines*.

You can get better performance for a table and minimize storage space by using the techniques listed here:

- Table Columns
- Row Format
- Indexes
- Joins
- Normalization

Table Columns

- Use the most efficient (smallest) data types possible. MySQL has many specialized types that save disk space and memory. For example, use the smaller integer types if possible to get smaller tables. `MEDIUMINT` is often a better choice than `INT` because a `MEDIUMINT` column uses 25% less space.
- Declare columns to be `NOT NULL` if possible. It makes SQL operations faster, by enabling better use of indexes and eliminating overhead for testing whether each value is `NULL`. You also save some storage space, one bit per column. If you really need `NULL` values in your tables, use them. Just avoid the default setting that allows `NULL` values in every column.

Row Format

- `InnoDB` tables are created using the `DYNAMIC` row format by default. To use a row format other than `DYNAMIC`, configure `innodb_default_row_format`, or specify the `ROW_FORMAT` option explicitly in a `CREATE TABLE` or `ALTER TABLE` statement.

  The compact family of row formats, which includes `COMPACT`, `DYNAMIC`, and `COMPRESSED`, decreases row storage space at the cost of increasing CPU use for some operations. If your workload is a typical one that is limited by cache hit rates and disk speed it is likely to be faster. If it is a rare case that is limited by CPU speed, it might be slower.

  The compact family of row formats also optimizes `CHAR` column storage when using a variable-length character set such as `utf8mb3` or `utf8mb4`. With `ROW_FORMAT=REDUNDANT`, `CHAR(N)` occupies *`N`* × the maximum byte length of the character set. Many languages can be written primarily using single-byte `utf8` characters, so a fixed storage length often wastes space. With the compact family of rows formats, `InnoDB` allocates a variable amount of storage in the range of *`N`* to *`N`* × the maximum byte length of the character set for these columns by stripping trailing spaces.

The minimum storage length is *N* bytes to facilitate in-place updates in typical cases. For more information, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

- To minimize space even further by storing table data in compressed form, specify `ROW_FORMAT=COMPRESSED` when creating `InnoDB` tables, or run the `myisampack` command on an existing `MyISAM` table. (`InnoDB` compressed tables are readable and writable, while `MyISAM` compressed tables are read-only.)
- For `MyISAM` tables, if you do not have any variable-length columns (`VARCHAR`, `TEXT`, or `BLOB` columns), a fixed-size row format is used. This is faster but may waste some space. See Section 16.2.3, “MyISAM Table Storage Formats”. You can hint that you want to have fixed length rows even if you have `VARCHAR` columns with the `CREATE TABLE` option `ROW_FORMAT=FIXED`.

Indexes

- The primary index of a table should be as short as possible. This makes identification of each row easy and efficient. For `InnoDB` tables, the primary key columns are duplicated in each secondary index entry, so a short primary key saves considerable space if you have many secondary indexes.
- Create only the indexes that you need to improve query performance. Indexes are good for retrieval, but slow down insert and update operations. If you access a table mostly by searching on a combination of columns, create a single composite index on them rather than a separate index for each column. The first part of the index should be the column most used. If you *always* use many columns when selecting from the table, the first column in the index should be the one with the most duplicates, to obtain better compression of the index.
- If it is very likely that a long string column has a unique prefix on the first number of characters, it is better to index only this prefix, using MySQL's support for creating an index on the leftmost part of the column (see Section 13.1.12, “CREATE INDEX Syntax”). Shorter indexes are faster, not only because they require less disk space, but because they also give you more hits in the index cache, and thus fewer disk seeks. See Section 5.1.1, “Configuring the Server”.

Joins

- In some circumstances, it can be beneficial to split into two a table that is scanned very often. This is especially true if it is a dynamic-format table and it is possible to use a smaller static format table that can be used to find the relevant rows when scanning the table.
- Declare columns with identical information in different tables with identical data types, to speed up joins based on the corresponding columns.
- Keep column names simple, so that you can use the same name across different tables and simplify join queries. For example, in a table named `customer`, use a column name of `name` instead of `customer_name`. To make your names portable to other SQL servers, consider keeping them shorter than 18 characters.

Normalization

- Normally, try to keep all data nonredundant (observing what is referred to in database theory as *third normal form*). Instead of repeating lengthy values such as names and addresses, assign them unique IDs, repeat these IDs as needed across multiple smaller tables, and join the tables in queries by referencing the IDs in the join clause.
- If speed is more important than disk space and the maintenance costs of keeping multiple copies of data, for example in a business intelligence scenario where you analyze all the data from large tables, you can relax the normalization rules, duplicating information or creating summary tables to gain more speed.

8.4.2 Optimizing MySQL Data Types

8.4.2.1 Optimizing for Numeric Data

- For unique IDs or other values that can be represented as either strings or numbers, prefer numeric columns to string columns. Since large numeric values can be stored in fewer bytes than the corresponding strings, it is faster and takes less memory to transfer and compare them.
- If you are using numeric data, it is faster in many cases to access information from a database (using a live connection) than to access a text file. Information in the database is likely to be stored in a more compact format than in the text file, so accessing it involves fewer disk accesses. You also save code in your application because you can avoid parsing the text file to find line and column boundaries.

8.4.2.2 Optimizing for Character and String Types

For character and string columns, follow these guidelines:

- Use binary collation order for fast comparison and sort operations, when you do not need language-specific collation features. You can use the `BINARY` operator to use binary collation within a particular query.
- When comparing values from different columns, declare those columns with the same character set and collation wherever possible, to avoid string conversions while running the query.
- For column values less than 8KB in size, use binary `VARCHAR` instead of `BLOB`. The `GROUP BY` and `ORDER BY` clauses can generate temporary tables, and these temporary tables can use the `MEMORY` storage engine if the original table does not contain any `BLOB` columns.
- If a table contains string columns such as name and address, but many queries do not retrieve those columns, consider splitting the string columns into a separate table and using join queries with a foreign key when necessary. When MySQL retrieves any value from a row, it reads a data block containing all the columns of that row (and possibly other adjacent rows). Keeping each row small, with only the most frequently used columns, allows more rows to fit in each data block. Such compact tables reduce disk I/O and memory usage for common queries.
- When you use a randomly generated value as a primary key in an `InnoDB` table, prefix it with an ascending value such as the current date and time if possible. When consecutive primary values are physically stored near each other, `InnoDB` can insert and retrieve them faster.
- See Section 8.4.2.1, “Optimizing for Numeric Data” for reasons why a numeric column is usually preferable to an equivalent string column.

8.4.2.3 Optimizing for BLOB Types

- When storing a large blob containing textual data, consider compressing it first. Do not use this technique when the entire table is compressed by `InnoDB` or `MyISAM`.
- For a table with several columns, to reduce memory requirements for queries that do not use the BLOB column, consider splitting the BLOB column into a separate table and referencing it with a join query when needed.
- Since the performance requirements to retrieve and display a BLOB value might be very different from other data types, you could put the BLOB-specific table on a different storage device or even a separate database instance. For example, to retrieve a BLOB might require a large sequential disk read that is better suited to a traditional hard drive than to an SSD device.
- See Section 8.4.2.2, “Optimizing for Character and String Types” for reasons why a binary `VARCHAR` column is sometimes preferable to an equivalent BLOB column.
- Rather than testing for equality against a very long text string, you can store a hash of the column value in a separate column, index that column, and test the hashed value in queries. (Use the `MD5()` or `CRC32()` function to produce the hash value.) Since hash functions can produce duplicate results for different inputs, you still include a clause `AND` *`blob_column = long_string_value`* in the query to guard against false matches; the performance benefit comes from the smaller, easily scanned index for the hashed values.

8.4.3 Optimizing for Many Tables

Some techniques for keeping individual queries fast involve splitting data across many tables. When the number of tables runs into the thousands or even millions, the overhead of dealing with all these tables becomes a new performance consideration.

8.4.3.1 How MySQL Opens and Closes Tables

When you execute a `mysqladmin status` command, you should see something like this:

```
Uptime: 426 Running threads: 1 Questions: 11082
Reloads: 1 Open tables: 12
```

The `Open tables` value of 12 can be somewhat puzzling if you have only six tables.

MySQL is multi-threaded, so there may be many clients issuing queries for a given table simultaneously. To minimize the problem with multiple client sessions having different states on the same table, the table is opened independently by each concurrent session. This uses additional memory but normally increases performance. With `MyISAM` tables, one extra file descriptor is required for the data file for each client that has the table open. (By contrast, the index file descriptor is shared between all sessions.)

The `table_open_cache` and `max_connections` system variables affect the maximum number of files the server keeps open. If you increase one or both of these values, you may run up against a limit imposed by your operating system on the per-process number of open file descriptors. Many operating systems permit you to increase the open-files limit, although the method varies widely from system to system. Consult your operating system documentation to determine whether it is possible to increase the limit and how to do so.

`table_open_cache` is related to `max_connections`. For example, for 200 concurrent running connections, specify a table cache size of at least `200 * N`, where *`N`* is the maximum number of tables per join in any of the queries which you execute. You must also reserve some extra file descriptors for temporary tables and files.

Make sure that your operating system can handle the number of open file descriptors implied by the `table_open_cache` setting. If `table_open_cache` is set too high, MySQL may run out of file descriptors and refuse connections, fail to perform queries, and be very unreliable.

You should also take into account the fact that the `MyISAM` storage engine needs two file descriptors for each unique open table. For a partitioned `MyISAM` table, two file descriptors are required for each partition of the opened table. (Note further that when `MyISAM` opens a partitioned table, it opens every partition of this table, whether or not a given partition is actually used. See MyISAM and partition file descriptor usage.) You can increase the number of file descriptors available to MySQL using the `--open-files-limit` startup option to `mysqld`. See Section B.5.2.17, “File Not Found and Similar Errors”.

The cache of open tables is kept at a level of `table_open_cache` entries. The server autosizes the cache size at startup. To set the size explicitly, set the `table_open_cache` system variable at startup. Note that MySQL may temporarily open more tables than this to execute queries.

MySQL closes an unused table and removes it from the table cache under the following circumstances:

- When the cache is full and a thread tries to open a table that is not in the cache.
- When the cache contains more than `table_open_cache` entries and a table in the cache is no longer being used by any threads.
- When a table flushing operation occurs. This happens when someone issues a `FLUSH TABLES` statement or executes a `mysqladmin flush-tables` or `mysqladmin refresh` command.

When the table cache fills up, the server uses the following procedure to locate a cache entry to use:

- Tables that are not currently in use are released, beginning with the table least recently used.
- If a new table needs to be opened, but the cache is full and no tables can be released, the cache is temporarily extended as necessary. When the cache is in a temporarily extended state and a table goes from a used to unused state, the table is closed and released from the cache.

A `MyISAM` table is opened for each concurrent access. This means the table needs to be opened twice if two threads access the same table or if a thread accesses the table twice in the same query (for example, by joining the table to itself). Each concurrent open requires an entry in the table cache. The first open of any `MyISAM` table takes two file descriptors: one for the data file and one for the index file. Each additional use of the table takes only one file descriptor for the data file. The index file descriptor is shared among all threads.

If you are opening a table with the `HANDLER tbl_name OPEN` statement, a dedicated table object is allocated for the thread. This table object is not shared by other threads and is not closed until the thread calls `HANDLER tbl_name CLOSE` or the thread terminates. When this happens, the table is put back in the table cache (if the cache is not full). See Section 13.2.4, “HANDLER Syntax”.

You can determine whether your table cache is too small by checking the `mysqld` status variable `Opened_tables`, which indicates the number of table-opening operations since the server started:

```
mysql> SHOW GLOBAL STATUS LIKE 'Opened_tables';
+---------------+-------+
| Variable_name | Value |
+---------------+-------+
| Opened_tables | 2741  |
+---------------+-------+
```

If the value is very large or increases rapidly, even when you have not issued many `FLUSH TABLES` statements, increase the table cache size. See Section 5.1.5, “Server System Variables”, and Section 5.1.7, “Server Status Variables”.

8.4.3.2 Disadvantages of Creating Many Tables in the Same Database

If you have many `MyISAM` tables in the same database directory, open, close, and create operations are slow. If you execute `SELECT` statements on many different tables, there is a little overhead when the table cache is full, because for every table that has to be opened, another must be closed. You can reduce this overhead by increasing the number of entries permitted in the table cache.

8.4.4 Internal Temporary Table Use in MySQL

In some cases, the server creates internal temporary tables while processing statements. Users have no direct control over when this occurs.

The server creates temporary tables under conditions such as these:

- Evaluation of `UNION` statements, with some exceptions described later.
- Evaluation of some views, such those that use the `TEMPTABLE` algorithm, `UNION`, or aggregation.
- Evaluation of derived tables (subqueries in the `FROM` clause).
- Evaluation of common table expressions (see Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”).
- Tables created for subquery or semi-join materialization (see Section 8.2.2, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions”).
- Evaluation of statements that contain an `ORDER BY` clause and a different `GROUP BY` clause, or for which the `ORDER BY` or `GROUP BY` contains columns from tables other than the first table in the join queue.

- Evaluation of `DISTINCT` combined with `ORDER BY` may require a temporary table.
- For queries that use the `SQL_SMALL_RESULT` modifier, MySQL uses an in-memory temporary table, unless the query also contains elements (described later) that require on-disk storage.
- Evaluation of multiple-table `UPDATE` statements.
- Evaluation of `GROUP_CONCAT()` or `COUNT(DISTINCT)` expressions.

To determine whether a statement requires a temporary table, use `EXPLAIN` and check the `Extra` column to see whether it says `Using temporary` (see Section 8.8.1, “Optimizing Queries with EXPLAIN”). `EXPLAIN` will not necessarily say `Using temporary` for derived or materialized temporary tables.

When the server creates an internal temporary table (either in memory or on disk), it increments the `Created_tmp_tables` status variable. If the server creates the table on disk (either initially or by converting an in-memory table) it increments the `Created_tmp_disk_tables` status variable.

Some query conditions prevent the use of an in-memory temporary table, in which case the server uses an on-disk table instead:

- Presence of a `BLOB` or `TEXT` column in the table
- Presence of any string column with a maximum length larger than 512 (bytes for binary strings, characters for nonbinary strings) in the `SELECT` list, if `UNION` or `UNION ALL` is used
- The `SHOW COLUMNS` and `DESCRIBE` statements use `BLOB` as the type for some columns, thus the temporary table used for the results is an on-disk table.

The server does not use a temporary table for `UNION` statements that meet certain qualifications. Instead, it retains from temporary table creation only the data structures necessary to perform result column typecasting. The table is not fully instantiated and no rows are written to or read from it; rows are sent directly to the client. The result is reduced memory and disk requirements, and smaller delay before the first row is sent to the client because the server need not wait until the last query block is executed. `EXPLAIN` and optimizer trace output reflects this execution strategy: The `UNION RESULT` query block is not present because that block corresponds to the part that reads from the temporary table.

These conditions qualify a `UNION` for evaluation without a temporary table:

- The union is `UNION ALL`, not `UNION` or `UNION DISTINCT`.
- There is no global `ORDER BY` clause.
- The union is not the top-level query block of an `{INSERT | REPLACE} ... SELECT ...` statement.
- Storage Engines Used for Temporary Tables
- Temporary Table Storage Format

Storage Engines Used for Temporary Tables

An internal temporary table can be held in memory and processed by the `TempTable` or `MEMORY` storage engine, or stored on disk by the `InnoDB` or `MyISAM` storage engine.

The `internal_tmp_mem_storage_engine` session variable, introduced in MySQL 8.0.2, defines the storage engine for in-memory internal temporary tables. Permitted values are `TempTable` (the default as of MySQL 8.0.2) and `MEMORY`.

The `TempTable` storage engine, introduced in MySQL 8.0.2, replaces the `MEMORY` storage engine as the default engine for in-memory internal temporary tables. The `TempTable` storage engine provides efficient storage for `VARCHAR` and `VARBINARY` columns.

The `temptable_max_ram` configuration option defines the maximum amount of memory that the `TempTable` storage engine can use before data is stored to disk. The default value is 1GiB.

The `memory/temptable/physical_ram` and `memory/temptable/physical_disk` Performance Schema instruments may be used to monitor `TempTable` memory allocation and disk storage. `memory/temptable/physical_ram` reports the amount of allocated RAM. `memory/temptable/physical_disk` reports the amount of allocated disk. If this instrument reports a value other than 0, the `temptable_max_ram` threshold was reached at some point. Data can be queried in Performance Schema memory summary tables such as `memory_summary_global_by_event_name`. See Section 25.11.15.10, “Memory Summary Tables”.

When using the `MEMORY` storage engine for in-memory temporary tables, MySQL automatically converts an in-memory temporary table to an on-disk table if it becomes too large. The maximum size for in-memory temporary tables is determined from whichever of the values of `tmp_table_size` and `max_heap_table_size` is smaller. This differs from `MEMORY` tables explicitly created with `CREATE TABLE`: For such tables, only the `max_heap_table_size` system variable determines how large the table is permitted to grow and there is no conversion to on-disk format.

The `tmp_table_size` and `max_heap_table_size` configuration options do not apply to the `TempTable` storage engine.

The `internal_tmp_disk_storage_engine` system variable determines which storage engine the server uses to manage on-disk internal temporary tables. Permitted values are `INNODB` (the default) and `MYISAM`.

For common table expressions (CTEs), the storage engine used for on-disk internal temporary tables cannot be `MyISAM`. If `internal_tmp_disk_storage_engine=MYISAM`, an error occurs for any attempt to materialize a CTE using an on-disk temporary table.

Note

Using `internal_tmp_disk_storage_engine=INNODB`, queries that generate temporary tables exceeding `InnoDB` row or column limits return `Row size too large` or `Too many columns` errors. The workaround is to set `internal_tmp_disk_storage_engine` to `MYISAM`.

Temporary Table Storage Format

When in-memory temporary tables are managed by the `TempTable` storage engine, rows that include `VARCHAR` and `VARBINARY` columns are represented in memory by an array of cells, with each cell containing a NULL flag, the data length, and a data pointer. Column values are placed in consecutive order after the array, in a single region of memory, without padding. Each cell in the array uses 16B of storage.

When in-memory temporary tables are managed by the `MEMORY` storage engine, fixed-length row format is used. `VARCHAR` and `VARBINARY` column values are padded to the maximum column length, in effect storing them as `CHAR` and `BINARY` columns.

On-disk temporary tables are managed by the `InnoDB` or `MyISAM` storage engine (depending on the `internal_tmp_disk_storage_engine` setting). Both engines store temporary tables using dynamic-width row format. Columns take only as much storage as needed, which reduces disk I/O and space requirements, and processing time compared to on-disk tables that use fixed-length rows.

When using the `MEMORY` storage engine, statements that initially create an internal temporary table in memory, and then convert it to an on-disk table, better performance might be achieved by skipping the conversion step and creating the table on disk to begin with.

The `big_tables` system variable can be used to force disk storage of internal temporary tables.

8.5 Optimizing for InnoDB Tables

`InnoDB` is the storage engine that MySQL customers typically use in production databases where reliability and concurrency are important. `InnoDB` is the default storage engine in MySQL. This section explains how to optimize database operations for `InnoDB` tables.

8.5.1 Optimizing Storage Layout for InnoDB Tables

- Once your data reaches a stable size, or a growing table has increased by tens or some hundreds of megabytes, consider using the `OPTIMIZE TABLE` statement to reorganize the table and compact any wasted space. The reorganized tables require less disk I/O to perform full table scans. This is a straightforward technique that can improve performance when other techniques such as improving index usage or tuning application code are not practical.

  `OPTIMIZE TABLE` copies the data part of the table and rebuilds the indexes. The benefits come from improved packing of data within indexes, and reduced fragmentation within the tablespaces and on disk. The benefits vary depending on the data in each table. You may find that there are significant gains for some and not for others, or that the gains decrease over time until you next optimize the table. This operation can be slow if the table is large or if the indexes being rebuilt do not fit into the buffer pool. The first run after adding a lot of data to a table is often much slower than later runs.

- In `InnoDB`, having a long `PRIMARY KEY` (either a single column with a lengthy value, or several columns that form a long composite value) wastes a lot of disk space. The primary key value for a row is duplicated in all the secondary index records that point to the same row. (See Section 15.8.2.1, “Clustered and Secondary Indexes”.) Create an `AUTO_INCREMENT` column as the primary key if your primary key is long, or index a prefix of a long `VARCHAR` column instead of the entire column.

- Use the `VARCHAR` data type instead of `CHAR` to store variable-length strings or for columns with many `NULL` values. A `CHAR(N)` column always takes *`N`* characters to store data, even if the string is shorter or its value is `NULL`. Smaller tables fit better in the buffer pool and reduce disk I/O.

  When using `COMPACT` row format (the default `InnoDB` format) and variable-length character sets, such as `utf8` or `sjis`, `CHAR(N)` columns occupy a variable amount of space, but still at least *`N`* bytes.

- For tables that are big, or contain lots of repetitive text or numeric data, consider using `COMPRESSED` row format. Less disk I/O is required to bring data into the buffer pool, or to perform full table scans. Before making a permanent decision, measure the amount of compression you can achieve by using `COMPRESSED` versus `COMPACT` row format.

8.5.2 Optimizing InnoDB Transaction Management

To optimize `InnoDB` transaction processing, find the ideal balance between the performance overhead of transactional features and the workload of your server. For example, an application might encounter performance issues if it commits thousands of times per second, and different performance issues if it commits only every 2-3 hours.

- The default MySQL setting `AUTOCOMMIT=1` can impose performance limitations on a busy database server. Where practical, wrap several related data change operations into a single transaction, by issuing `SET AUTOCOMMIT=0` or a `START TRANSACTION` statement, followed by a `COMMIT` statement after making all the changes.

  `InnoDB` must flush the log to disk at each transaction commit if that transaction made modifications to the database. When each change is followed by a commit (as with the default autocommit setting), the I/O throughput of the storage device puts a cap on the number of potential operations per second.

- Alternatively, for transactions that consist only of a single `SELECT` statement, turning on `AUTOCOMMIT` helps `InnoDB` to recognize read-only transactions and optimize them. See Section 8.5.3, “Optimizing InnoDB Read-Only Transactions” for requirements.

- Avoid performing rollbacks after inserting, updating, or deleting huge numbers of rows. If a big transaction is slowing down server performance, rolling it back can make the problem worse, potentially taking several times as long to perform as the original data change operations. Killing the database process does not help, because the rollback starts again on server startup.

  To minimize the chance of this issue occurring:

  - Increase the size of the buffer pool so that all the data change changes can be cached rather than immediately written to disk.
  - Set `innodb_change_buffering=all` so that update and delete operations are buffered in addition to inserts.
  - Consider issuing `COMMIT` statements periodically during the big data change operation, possibly breaking a single delete or update into multiple statements that operate on smaller numbers of rows.

  To get rid of a runaway rollback once it occurs, increase the buffer pool so that the rollback becomes CPU-bound and runs fast, or kill the server and restart with `innodb_force_recovery=3`, as explained in Section 15.17.2, “InnoDB Recovery”.

  This issue is expected to be infrequent with the default setting `innodb_change_buffering=all`, which allows update and delete operations to be cached in memory, making them faster to perform in the first place, and also faster to roll back if needed. Make sure to use this parameter setting on servers that process long-running transactions with many inserts, updates, or deletes.

- If you can afford the loss of some of the latest committed transactions if a crash occurs, you can set the `innodb_flush_log_at_trx_commit` parameter to 0. `InnoDB` tries to flush the log once per second anyway, although the flush is not guaranteed.

- When rows are modified or deleted, the rows and associated undo logs are not physically removed immediately, or even immediately after the transaction commits. The old data is preserved until transactions that started earlier or concurrently are finished, so that those transactions can access the previous state of modified or deleted rows. Thus, a long-running transaction can prevent `InnoDB` from purging data that was changed by a different transaction.

- When rows are modified or deleted within a long-running transaction, other transactions using the `READ COMMITTED` and `REPEATABLE READ` isolation levels have to do more work to reconstruct the older data if they read those same rows.

- When a long-running transaction modifies a table, queries against that table from other transactions do not make use of the covering index technique. Queries that normally could retrieve all the result columns from a secondary index, instead look up the appropriate values from the table data.

  If secondary index pages are found to have a `PAGE_MAX_TRX_ID` that is too new, or if records in the secondary index are delete-marked, `InnoDB` may need to look up records using a clustered index.

8.5.3 Optimizing InnoDB Read-Only Transactions

`InnoDB` can avoid the overhead associated with setting up the transaction ID (`TRX_ID` field) for transactions that are known to be read-only. A transaction ID is only needed for a transaction that might perform write operations or locking reads such as `SELECT ... FOR UPDATE`. Eliminating unnecessary transaction IDs reduces the size of internal data structures that are consulted each time a query or data change statement constructs a read view.

`InnoDB` detects read-only transactions when:

- The transaction is started with the `START TRANSACTION READ ONLY` statement. In this case, attempting to make changes to the database (for `InnoDB`, `MyISAM`, or other types of tables) causes an error, and the transaction continues in read-only state:

```
ERROR 1792 (25006): Cannot execute statement in a READ ONLY transaction.
```

You can still make changes to session-specific temporary tables in a read-only transaction, or issue locking queries for them, because those changes and locks are not visible to any other transaction.

- The `autocommit` setting is turned on, so that the transaction is guaranteed to be a single statement, and the single statement making up the transaction is a "non-locking" `SELECT` statement. That is, a `SELECT` that does not use a `FOR UPDATE` or `LOCK IN SHARED MODE` clause.
- The transaction is started without the `READ ONLY` option, but no updates or statements that explicitly lock rows have been executed yet. Until updates or explicit locks are required, a transaction stays in read-only mode.

Thus, for a read-intensive application such as a report generator, you can tune a sequence of `InnoDB` queries by grouping them inside `START TRANSACTION READ ONLY` and `COMMIT`, or by turning on the `autocommit` setting before running the `SELECT` statements, or simply by avoiding any data change statements interspersed with the queries.

For information about `START TRANSACTION` and `autocommit`, see Section 13.3.1, "START TRANSACTION, COMMIT, and ROLLBACK Syntax".

Note

Transactions that qualify as auto-commit, non-locking, and read-only (AC-NL-RO) are kept out of certain internal `InnoDB` data structures and are therefore not listed in `SHOW ENGINE INNODB STATUS` output.

8.5.4 Optimizing InnoDB Redo Logging

Consider the following guidelines for optimizing redo logging:

- Make your redo log files big, even as big as the buffer pool. When `InnoDB` has written the redo log files full, it must write the modified contents of the buffer pool to disk in a checkpoint. Small redo log files cause many unnecessary disk writes. Although historically big redo log files caused lengthy recovery times, recovery is now much faster and you can confidently use large redo log files.

  The size and number of redo log files are configured using the `innodb_log_file_size` and `innodb_log_files_in_group` configuration options. For information about modifying an existing redo log file configuration, see Section 15.7.2, "Changing the Number or Size of InnoDB Redo Log Files".

- Consider increasing the size of the log buffer. A large log buffer enables large transactions to run without a need to write the log to disk before the transactions commit. Thus, if you have transactions that update, insert, or delete many rows, making the log buffer larger saves disk I/O. Log buffer size is configured using the `innodb_log_buffer_size` configuration option.

8.5.5 Bulk Data Loading for InnoDB Tables

These performance tips supplement the general guidelines for fast inserts in Section 8.2.5.1, "Optimizing INSERT Statements".

- When importing data into `InnoDB`, turn off autocommit mode, because it performs a log flush to disk for every insert. To disable autocommit during your import operation, surround it with `SET autocommit` and `COMMIT` statements:

```
SET autocommit=0;
... SQL import statements ...
COMMIT;
```

The `mysqldump` option `--opt` creates dump files that are fast to import into an `InnoDB` table, even without wrapping them with the `SET autocommit` and `COMMIT` statements.

- If you have `UNIQUE` constraints on secondary keys, you can speed up table imports by temporarily turning off the uniqueness checks during the import session:

```
SET unique_checks=0;
... SQL import statements ...
SET unique_checks=1;
```

For big tables, this saves a lot of disk I/O because `InnoDB` can use its change buffer to write secondary index records in a batch. Be certain that the data contains no duplicate keys.

- If you have `FOREIGN KEY` constraints in your tables, you can speed up table imports by turning off the foreign key checks for the duration of the import session:

```
SET foreign_key_checks=0;
... SQL import statements ...
SET foreign_key_checks=1;
```

For big tables, this can save a lot of disk I/O.

- Use the multiple-row `INSERT` syntax to reduce communication overhead between the client and the server if you need to insert many rows:

```
INSERT INTO yourtable VALUES (1,2), (5,5), ...;
```

This tip is valid for inserts into any table, not just `InnoDB` tables.

- When doing bulk inserts into tables with auto-increment columns, set `innodb_autoinc_lock_mode` to 2 instead of the default value 1. See Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB” for details.

- When performing bulk inserts, it is faster to insert rows in `PRIMARY KEY` order. `InnoDB` tables use a clustered index, which makes it relatively fast to use data in the order of the `PRIMARY KEY`. Performing bulk inserts in `PRIMARY KEY` order is particularly important for tables that do not fit entirely within the buffer pool.

- For optimal performance when loading data into an `InnoDB` `FULLTEXT` index, follow this set of steps:

  1. Define a column `FTS_DOC_ID` at table creation time, of type `BIGINT UNSIGNED NOT NULL`, with a unique index named `FTS_DOC_ID_INDEX`. For example:

```
CREATE TABLE t1 (
FTS_DOC_ID BIGINT unsigned NOT NULL AUTO_INCREMENT,
title varchar(255) NOT NULL DEFAULT '',
text mediumtext NOT NULL,
PRIMARY KEY (`FTS_DOC_ID`)
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4;
CREATE UNIQUE INDEX FTS_DOC_ID_INDEX on t1(FTS_DOC_ID);
```

  2. Load the data into the table.

  3. Create the `FULLTEXT` index after the data is loaded.

Note

When adding `FTS_DOC_ID` column at table creation time, ensure that the `FTS_DOC_ID` column is updated when the `FULLTEXT` indexed column

> is updated, as the `FTS_DOC_ID` must increase monotonically with each `INSERT` or `UPDATE`. If you choose not to add the `FTS_DOC_ID` at table creation time and have `InnoDB` manage DOC IDs for you, `InnoDB` will add the `FTS_DOC_ID` as a hidden column with the next `CREATE FULLTEXT INDEX` call. This approach, however, requires a table rebuild which will impact performance.

8.5.6 Optimizing InnoDB Queries

To tune queries for `InnoDB` tables, create an appropriate set of indexes on each table. See Section 8.3.1, “How MySQL Uses Indexes” for details. Follow these guidelines for `InnoDB` indexes:

- Because each `InnoDB` table has a primary key (whether you request one or not), specify a set of primary key columns for each table, columns that are used in the most important and time-critical queries.
- Do not specify too many or too long columns in the primary key, because these column values are duplicated in each secondary index. When an index contains unnecessary data, the I/O to read this data and memory to cache it reduce the performance and scalability of the server.
- Do not create a separate secondary index for each column, because each query can only make use of one index. Indexes on rarely tested columns or columns with only a few different values might not be helpful for any queries. If you have many queries for the same table, testing different combinations of columns, try to create a small number of concatenated indexes rather than a large number of single-column indexes. If an index contains all the columns needed for the result set (known as a covering index), the query might be able to avoid reading the table data at all.
- If an indexed column cannot contain any `NULL` values, declare it as `NOT NULL` when you create the table. The optimizer can better determine which index is most effective to use for a query, when it knows whether each column contains `NULL` values.
- You can optimize single-query transactions for `InnoDB` tables, using the technique in Section 8.5.3, “Optimizing InnoDB Read-Only Transactions”.
- If you often have recurring queries for tables that are not updated frequently, enable the query cache:

```
[mysqld]
query_cache_type = 1
query_cache_size = 10M
```

8.5.7 Optimizing InnoDB DDL Operations

- For DDL operations on tables and indexes (`CREATE`, `ALTER`, and `DROP` statements), the most significant aspect for `InnoDB` tables is that creating and dropping secondary indexes is much faster in MySQL 5.5 and higher, than in earlier releases. See Section 15.12.1, “Online DDL Overview” for details.
- “Fast index creation” makes it faster in some cases to drop an index before loading data into a table, then re-create the index after loading the data.
- Use `TRUNCATE TABLE` to empty a table, not `DELETE FROM` *`tbl_name`*. Foreign key constraints can make a `TRUNCATE` statement work like a regular `DELETE` statement, in which case a sequence of commands like `DROP TABLE` and `CREATE TABLE` might be fastest.
- Because the primary key is integral to the storage layout of each `InnoDB` table, and changing the definition of the primary key involves reorganizing the whole table, always set up the primary key as part of the `CREATE TABLE` statement, and plan ahead so that you do not need to `ALTER` or `DROP` the primary key afterward.

8.5.8 Optimizing InnoDB Disk I/O

If you follow best practices for database design and tuning techniques for SQL operations, but your database is still slow due to heavy disk I/O activity, consider these disk I/O optimizations. If the Unix `top` tool or the Windows Task Manager shows that the CPU usage percentage with your workload is less than 70%, your workload is probably disk-bound.

- Increase buffer pool size

  When table data is cached in the `InnoDB` buffer pool, it can be accessed repeatedly by queries without requiring any disk I/O. Specify the size of the buffer pool with the `innodb_buffer_pool_size` option. This memory area is important enough that it is typically recommended that `innodb_buffer_pool_size` is configured to 50 to 75 percent of system memory. For more information see, Section 8.12.3.1, “How MySQL Uses Memory”.

- Adjust the flush method

  In some versions of GNU/Linux and Unix, flushing files to disk with the Unix `fsync()` call (which `InnoDB` uses by default) and similar methods is surprisingly slow. If database write performance is an issue, conduct benchmarks with the `innodb_flush_method` parameter set to `O_DSYNC`.

- Use a noop or deadline I/O scheduler with native AIO on Linux

  `InnoDB` uses the asynchronous I/O subsystem (native AIO) on Linux to perform read-ahead and write requests for data file pages. This behavior is controlled by the `innodb_use_native_aio` configuration option, which is enabled by default. With native AIO, the type of I/O scheduler has greater influence on I/O performance. Generally, noop and deadline I/O schedulers are recommended. Conduct benchmarks to determine which I/O scheduler provides the best results for your workload and environment. For more information, see Section 15.6.7, “Using Asynchronous I/O on Linux”.

- Use direct I/O on Solaris 10 for x86_64 architecture

  When using the `InnoDB` storage engine on Solaris 10 for x86_64 architecture (AMD Opteron), use direct I/O for `InnoDB`-related files to avoid degradation of `InnoDB` performance. To use direct I/O for an entire UFS file system used for storing `InnoDB`-related files, mount it with the `forcedirectio` option; see `mount_ufs(1M)`. (The default on Solaris 10/x86_64 is *not* to use this option.) To apply direct I/O only to `InnoDB` file operations rather than the whole file system, set `innodb_flush_method = O_DIRECT`. With this setting, `InnoDB` calls `directio()` instead of `fcntl()` for I/O to data files (not for I/O to log files).

- Use raw storage for data and log files with Solaris 2.6 or later

  When using the `InnoDB` storage engine with a large `innodb_buffer_pool_size` value on any release of Solaris 2.6 and up and any platform (sparc/x86/x64/amd64), conduct benchmarks with `InnoDB` data files and log files on raw devices or on a separate direct I/O UFS file system, using the `forcedirectio` mount option as described earlier. (It is necessary to use the mount option rather than setting `innodb_flush_method` if you want direct I/O for the log files.) Users of the Veritas file system VxFS should use the `convosync=direct` mount option.

  Do not place other MySQL data files, such as those for `MyISAM` tables, on a direct I/O file system. Executables or libraries *must not* be placed on a direct I/O file system.

- Use additional storage devices

  Additional storage devices could be used to set up a RAID configuration. For related information, see Section 8.12.1, “Optimizing Disk I/O”.

  Alternatively, `InnoDB` tablespace data files and log files can be placed on different physical disks. For more information, refer to the following sections:

  - Section 15.6.1, “InnoDB Startup Configuration”

- Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory”
- Creating a General Tablespace
- Section 15.8.1.3, “Moving or Copying InnoDB Tables”

- Consider non-rotational storage

  Non-rotational storage generally provides better performance for random I/O operations; and rotational storage for sequential I/O operations. When distributing data and log files across rotational and non-rotational storage devices, consider the type of I/O operations that are predominantly performed on each file.

  Random I/O-oriented files typically include file-per-table and general tablespace data files, undo tablespace files, and temporary tablespace files. Sequential I/O-oriented files include `InnoDB` system tablespace files (due to doublewrite buffering and change buffering) and log files such as binary log files and redo log files.

  Review settings for the following configuration options when using non-rotational storage:

  - `innodb_checksum_algorithm`

    The `crc32` option uses a faster checksum algorithm and is recommended for fast storage systems.

  - `innodb_flush_neighbors`

    This option optimizes I/O for rotational storage devices. Disable it for non-rotational storage or a mix of rotational and non-rotational storage.

  - `innodb_io_capacity`

    The default setting of 200 is generally sufficient for a lower-end non-rotational storage device. For higher-end, bus-attached devices, consider a higher setting such as 1000.

  - `innodb_io_capacity_max`

    The default value of 2000 is intended for workloads that use non-rotational storage. For a high-end, bus-attached non-rotational storage device, consider a higher setting such as 2500.

  - `innodb_log_compressed_pages`

    If redo logs are on non-rotational storage, consider disabling this option to reduce logging. See Disable logging of compressed pages.

  - `innodb_log_file_size`

    If redo logs are on non-rotational storage, configure this option to maximize caching and write combining.

  - `innodb_page_size`

    Consider using a page size that matches the internal sector size of the disk. Early-generation SSD devices often have a 4k sector size. Some newer devices have a 16k sector size. The default `InnoDB` page size is 16k. Keeping the page size close to the storage device block size minimizes the amount of unchanged data that is rewritten to disk.

  - `binlog_row_image`

    If binary logs are on non-rotational storage and all tables have primary keys, consider setting this option to `minimal` to reduce logging.

Ensure that TRIM support is enabled for your operating system. It is typically enabled by default.

- Increase I/O capacity to avoid backlogs

  If throughput drops periodically because of `InnoDB` checkpoint operations, consider increasing the value of the `innodb_io_capacity` configuration option. Higher values cause more frequent flushing, avoiding the backlog of work that can cause dips in throughput.

- Lower I/O capacity if flushing does not fall behind

  If the system is not falling behind with `InnoDB` flushing operations, consider lowering the value of the `innodb_io_capacity` configuration option. Typically, you keep this option value as low as practical, but not so low that it causes periodic drops in throughput as mentioned in the preceding bullet. In a typical scenario where you could lower the option value, you might see a combination like this in the output from `SHOW ENGINE INNODB STATUS`:

  - History list length low, below a few thousand.
  - Insert buffer merges close to rows inserted.
  - Modified pages in buffer pool consistently well below `innodb_max_dirty_pages_pct` of the buffer pool. (Measure at a time when the server is not doing bulk inserts; it is normal during bulk inserts for the modified pages percentage to rise significantly.)
  - `Log sequence number - Last checkpoint` is at less than 7/8 or ideally less than 6/8 of the total size of the `InnoDB` log files.

- Store system tablespace files on Fusion-io devices

  You can take advantage of a doublewrite buffer-related I/O optimization by storing system tablespace files (“ibdata files”) on Fusion-io devices that support atomic writes. In this case, doublewrite buffering (`innodb_doublewrite`) is automatically disabled and Fusion-io atomic writes are used for all data files. This feature is only supported on Fusion-io hardware and is only enabled for Fusion-io NVMFS on Linux. To take full advantage of this feature, an `innodb_flush_method` setting of `O_DIRECT` is recommended.

  Note

  Because the doublewrite buffer setting is global, doublewrite buffering is also disabled for data files residing on non-Fusion-io hardware.

- Disable logging of compressed pages

  When using the `InnoDB` table compression feature, images of re-compressed pages are written to the redo log when changes are made to compressed data. This behavior is controlled by `innodb_log_compressed_pages`, which is enabled by default to prevent corruption that can occur if a different version of the `zlib` compression algorithm is used during recovery. If you are certain that the `zlib` version will not change, disable `innodb_log_compressed_pages` to reduce redo log generation for workloads that modify compressed data.

8.5.9 Optimizing InnoDB Configuration Variables

Different settings work best for servers with light, predictable loads, versus servers that are running near full capacity all the time, or that experience spikes of high activity.

Because the `InnoDB` storage engine performs many of its optimizations automatically, many performance-tuning tasks involve monitoring to ensure that the database is performing well, and changing configuration options when performance drops. See Section 15.15, “InnoDB Integration with MySQL Performance Schema” for information about detailed `InnoDB` performance monitoring.

The main configuration steps you can perform include:

- Controlling the types of data change operations for which `InnoDB` buffers the changed data, to avoid frequent small disk writes. See Section 15.6.4, “Configuring InnoDB Change Buffering”. Because the default is to buffer all types of data change operations, only change this setting if you need to reduce the amount of buffering.
- Turning the adaptive hash indexing feature on and off using the `innodb_adaptive_hash_index` option. See Section 15.4.3, “Adaptive Hash Index” for more information. You might change this setting during periods of unusual activity, then restore it to its original setting.
- Setting a limit on the number of concurrent threads that `InnoDB` processes, if context switching is a bottleneck. See Section 15.6.5, “Configuring Thread Concurrency for InnoDB”.
- Controlling the amount of prefetching that `InnoDB` does with its read-ahead operations. When the system has unused I/O capacity, more read-ahead can improve the performance of queries. Too much read-ahead can cause periodic drops in performance on a heavily loaded system. See Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”.
- Increasing the number of background threads for read or write operations, if you have a high-end I/O subsystem that is not fully utilized by the default values. See Section 15.6.6, “Configuring the Number of Background InnoDB I/O Threads”.
- Controlling how much I/O `InnoDB` performs in the background. See Section 15.6.8, “Configuring the InnoDB Master Thread I/O Rate”. You might scale back this setting if you observe periodic drops in performance.
- Controlling the algorithm that determines when `InnoDB` performs certain types of background writes. See Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing”. The algorithm works for some types of workloads but not others, so might turn off this setting if you observe periodic drops in performance.
- Taking advantage of multicore processors and their cache memory configuration, to minimize delays in context switching. See Section 15.6.9, “Configuring Spin Lock Polling”.
- Preventing one-time operations such as table scans from interfering with the frequently accessed data stored in the `InnoDB` buffer cache. See Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”.
- Adjusting log files to a size that makes sense for reliability and crash recovery. `InnoDB` log files have often been kept small to avoid long startup times after a crash. Optimizations introduced in MySQL 5.5 speed up certain steps of the crash recovery process. In particular, scanning the redo log and applying the redo log are faster due to improved algorithms for memory management. If you have kept your log files artificially small to avoid long startup times, you can now consider increasing log file size to reduce the I/O that occurs due recycling of redo log records.
- Configuring the size and number of instances for the `InnoDB` buffer pool, especially important for systems with multi-gigabyte buffer pools. See Section 15.6.3.3, “Configuring Multiple Buffer Pool Instances”.
- Increasing the maximum number of concurrent transactions, which dramatically improves scalability for the busiest databases. See Section 15.4.7, “Undo Logs”.
- Moving purge operations (a type of garbage collection) into a background thread. See Section 15.6.10, “Configuring InnoDB Purge Scheduling”. To effectively measure the results of this setting, tune the other I/O-related and thread-related configuration settings first.
- Reducing the amount of switching that `InnoDB` does between concurrent threads, so that SQL operations on a busy server do not queue up and form a “traffic jam”. Set a value for the `innodb_thread_concurrency` option, up to approximately 32 for a high-powered modern system. Increase the value for the `innodb_concurrency_tickets` option, typically to 5000 or so. This combination of options sets a cap on the number of threads that `InnoDB` processes at any one time, and allows each thread to do substantial work before being swapped out, so that the number of waiting threads stays low and operations can complete without excessive context switching.

8.5.10 Optimizing InnoDB for Systems with Many Tables

- If you have configured non-persistent optimizer statistics (a non-default configuration), `InnoDB` computes index cardinality values for a table the first time that table is accessed after startup, instead of storing such values in the table. This step can take significant time on systems that partition the data into many tables. Since this overhead only applies to the initial table open operation, to “warm up” a table for later use, access it immediately after startup by issuing a statement such as `SELECT 1 FROM tbl_name LIMIT 1`.

  Optimizer statistics are persisted to disk by default, enabled by the `innodb_stats_persistent` configuration option. For information about persistent optimizer statistics, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

8.6 Optimizing for MyISAM Tables

The `MyISAM` storage engine performs best with read-mostly data or with low-concurrency operations, because table locks limit the ability to perform simultaneous updates. In MySQL, `InnoDB` is the default storage engine rather than `MyISAM`.

8.6.1 Optimizing MyISAM Queries

Some general tips for speeding up queries on `MyISAM` tables:

- To help MySQL better optimize queries, use `ANALYZE TABLE` or run `myisamchk --analyze` on a table after it has been loaded with data. This updates a value for each index part that indicates the average number of rows that have the same value. (For unique indexes, this is always 1.) MySQL uses this to decide which index to choose when you join two tables based on a nonconstant expression. You can check the result from the table analysis by using `SHOW INDEX FROM tbl_name` and examining the `Cardinality` value. `myisamchk --description --verbose` shows index distribution information.

- To sort an index and data according to an index, use `myisamchk --sort-index --sort-records=1` (assuming that you want to sort on index 1). This is a good way to make queries faster if you have a unique index from which you want to read all rows in order according to the index. The first time you sort a large table this way, it may take a long time.

- Try to avoid complex `SELECT` queries on `MyISAM` tables that are updated frequently, to avoid problems with table locking that occur due to contention between readers and writers.

- `MyISAM` supports concurrent inserts: If a table has no free blocks in the middle of the data file, you can `INSERT` new rows into it at the same time that other threads are reading from the table. If it is important to be able to do this, consider using the table in ways that avoid deleting rows. Another possibility is to run `OPTIMIZE TABLE` to defragment the table after you have deleted a lot of rows from it. This behavior is altered by setting the `concurrent_insert` variable. You can force new rows to be appended (and therefore permit concurrent inserts), even in tables that have deleted rows. See Section 8.11.3, “Concurrent Inserts”.

- For `MyISAM` tables that change frequently, try to avoid all variable-length columns (`VARCHAR`, `BLOB`, and `TEXT`). The table uses dynamic row format if it includes even a single variable-length column. See Chapter 16, *Alternative Storage Engines*.

- It is normally not useful to split a table into different tables just because the rows become large. In accessing a row, the biggest performance hit is the disk seek needed to find the first byte of the row. After finding the data, most modern disks can read the entire row fast enough for most applications. The only cases where splitting up a table makes an appreciable difference is if it is a `MyISAM` table using dynamic row format that you can change to a fixed row size, or if you very often need to scan the table but do not need most of the columns. See Chapter 16, *Alternative Storage Engines*.

- Use `ALTER TABLE ... ORDER BY expr1, expr2, ...` if you usually retrieve rows in *`expr1, expr2, ...`* order. By using this option after extensive changes to the table, you may be able to get higher performance.
- If you often need to calculate results such as counts based on information from a lot of rows, it may be preferable to introduce a new table and update the counter in real time. An update of the following form is very fast:

  ```
  UPDATE tbl_name SET count_col=count_col+1 WHERE key_col=constant;
  ```

  This is very important when you use MySQL storage engines such as `MyISAM` that has only table-level locking (multiple readers with single writers). This also gives better performance with most database systems, because the row locking manager in this case has less to do.
- Use `OPTIMIZE TABLE` periodically to avoid fragmentation with dynamic-format `MyISAM` tables. See Section 16.2.3, “MyISAM Table Storage Formats”.
- Declaring a `MyISAM` table with the `DELAY_KEY_WRITE=1` table option makes index updates faster because they are not flushed to disk until the table is closed. The downside is that if something kills the server while such a table is open, you must ensure that the table is okay by running the server with the `--myisam-recover-options` option, or by running `myisamchk` before restarting the server. (However, even in this case, you should not lose anything by using `DELAY_KEY_WRITE`, because the key information can always be generated from the data rows.)
- Strings are automatically prefix- and end-space compressed in `MyISAM` indexes. See Section 13.1.12, “CREATE INDEX Syntax”.
- You can increase performance by caching queries or answers in your application and then executing many inserts or updates together. Locking the table during this operation ensures that the index cache is only flushed once after all updates. You can also take advantage of MySQL's query cache to achieve similar results; see Section 8.10.3, “The MySQL Query Cache”.

8.6.2 Bulk Data Loading for MyISAM Tables

These performance tips supplement the general guidelines for fast inserts in Section 8.2.5.1, “Optimizing INSERT Statements”.

- For a `MyISAM` table, you can use concurrent inserts to add rows at the same time that `SELECT` statements are running, if there are no deleted rows in middle of the data file. See Section 8.11.3, “Concurrent Inserts”.
- With some extra work, it is possible to make `LOAD DATA INFILE` run even faster for a `MyISAM` table when the table has many indexes. Use the following procedure:

  1. Execute a `FLUSH TABLES` statement or a `mysqladmin flush-tables` command.
  2. Use `myisamchk --keys-used=0 -rq /path/to/db/tbl_name` to remove all use of indexes for the table.
  3. Insert data into the table with `LOAD DATA INFILE`. This does not update any indexes and therefore is very fast.
  4. If you intend only to read from the table in the future, use `myisampack` to compress it. See Section 16.2.3.3, “Compressed Table Characteristics”.
  5. Re-create the indexes with `myisamchk -rq /path/to/db/tbl_name`. This creates the index tree in memory before writing it to disk, which is much faster than updating the index during `LOAD DATA INFILE` because it avoids lots of disk seeks. The resulting index tree is also perfectly balanced.
  6. Execute a `FLUSH TABLES` statement or a `mysqladmin flush-tables` command.

`LOAD DATA INFILE` performs the preceding optimization automatically if the `MyISAM` table into which you insert data is empty. The main difference between automatic optimization and using the procedure explicitly is that you can let `myisamchk` allocate much more temporary memory for the index creation than you might want the server to allocate for index re-creation when it executes the `LOAD DATA INFILE` statement.

You can also disable or enable the nonunique indexes for a `MyISAM` table by using the following statements rather than `myisamchk`. If you use these statements, you can skip the `FLUSH TABLE` operations:

```
ALTER TABLE tbl_name DISABLE KEYS;
ALTER TABLE tbl_name ENABLE KEYS;
```

- To speed up `INSERT` operations that are performed with multiple statements for nontransactional tables, lock your tables:

```
LOCK TABLES a WRITE;
INSERT INTO a VALUES (1,23),(2,34),(4,33);
INSERT INTO a VALUES (8,26),(6,29);
...
UNLOCK TABLES;
```

This benefits performance because the index buffer is flushed to disk only once, after all `INSERT` statements have completed. Normally, there would be as many index buffer flushes as there are `INSERT` statements. Explicit locking statements are not needed if you can insert all rows with a single `INSERT`.

Locking also lowers the total time for multiple-connection tests, although the maximum wait time for individual connections might go up because they wait for locks. Suppose that five clients attempt to perform inserts simultaneously as follows:

- Connection 1 does 1000 inserts
- Connections 2, 3, and 4 do 1 insert
- Connection 5 does 1000 inserts

If you do not use locking, connections 2, 3, and 4 finish before 1 and 5. If you use locking, connections 2, 3, and 4 probably do not finish before 1 or 5, but the total time should be about 40% faster.

`INSERT`, `UPDATE`, and `DELETE` operations are very fast in MySQL, but you can obtain better overall performance by adding locks around everything that does more than about five successive inserts or updates. If you do very many successive inserts, you could do a `LOCK TABLES` followed by an `UNLOCK TABLES` once in a while (each 1,000 rows or so) to permit other threads to access table. This would still result in a nice performance gain.

`INSERT` is still much slower for loading data than `LOAD DATA INFILE`, even when using the strategies just outlined.

- To increase performance for `MyISAM` tables, for both `LOAD DATA INFILE` and `INSERT`, enlarge the key cache by increasing the `key_buffer_size` system variable. See Section 5.1.1, “Configuring the Server”.

8.6.3 Optimizing REPAIR TABLE Statements

`REPAIR TABLE` for `MyISAM` tables is similar to using `myisamchk` for repair operations, and some of the same performance optimizations apply:

- `myisamchk` has variables that control memory allocation. You may be able to its improve performance by setting these variables, as described in Section 4.6.4.6, “myisamchk Memory Usage”.
- For `REPAIR TABLE`, the same principle applies, but because the repair is done by the server, you set server system variables instead of `myisamchk` variables. Also, in addition to setting memory-allocation variables, increasing the `myisam_max_sort_file_size` system variable increases the likelihood that the repair will use the faster filesort method and avoid the slower repair by key cache method. Set the variable to the maximum file size for your system, after checking to be sure that there is enough free space to hold a copy of the table files. The free space must be available in the file system containing the original table files.

Suppose that a `myisamchk` table-repair operation is done using the following options to set its memory-allocation variables:

```
--key_buffer_size=128M --myisam_sort_buffer_size=256M
--read_buffer_size=64M --write_buffer_size=64M
```

Some of those `myisamchk` variables correspond to server system variables:

`myisamchk` Variable	**System Variable**
`key_buffer_size`	`key_buffer_size`
`myisam_sort_buffer_size`	`myisam_sort_buffer_size`
`read_buffer_size`	`read_buffer_size`
`write_buffer_size`	none

Each of the server system variables can be set at runtime, and some of them (`myisam_sort_buffer_size`, `read_buffer_size`) have a session value in addition to a global value. Setting a session value limits the effect of the change to your current session and does not affect other users. Changing a global-only variable (`key_buffer_size`, `myisam_max_sort_file_size`) affects other users as well. For `key_buffer_size`, you must take into account that the buffer is shared with those users. For example, if you set the `myisamchk key_buffer_size` variable to 128MB, you could set the corresponding `key_buffer_size` system variable larger than that (if it is not already set larger), to permit key buffer use by activity in other sessions. However, changing the global key buffer size invalidates the buffer, causing increased disk I/O and slowdown for other sessions. An alternative that avoids this problem is to use a separate key cache, assign to it the indexes from the table to be repaired, and deallocate it when the repair is complete. See Section 8.10.2.2, “Multiple Key Caches”.

Based on the preceding remarks, a `REPAIR TABLE` operation can be done as follows to use settings similar to the `myisamchk` command. Here a separate 128MB key buffer is allocated and the file system is assumed to permit a file size of at least 100GB.

```
SET SESSION myisam_sort_buffer_size = 256*1024*1024;
SET SESSION read_buffer_size = 64*1024*1024;
SET GLOBAL myisam_max_sort_file_size = 100*1024*1024*1024;
SET GLOBAL repair_cache.key_buffer_size = 128*1024*1024;
CACHE INDEX tbl_name IN repair_cache;
LOAD INDEX INTO CACHE tbl_name;
REPAIR TABLE tbl_name ;
SET GLOBAL repair_cache.key_buffer_size = 0;
```

If you intend to change a global variable but want to do so only for the duration of a `REPAIR TABLE` operation to minimally affect other users, save its value in a user variable and restore it afterward. For example:

```
SET @old_myisam_sort_buffer_size = @@global.myisam_max_sort_file_size;
SET GLOBAL myisam_max_sort_file_size = 100*1024*1024*1024;
REPAIR TABLE tbl_name ;
SET GLOBAL myisam_max_sort_file_size = @old_myisam_max_sort_file_size;
```

The system variables that affect `REPAIR TABLE` can be set globally at server startup if you want the values to be in effect by default. For example, add these lines to the server `my.cnf` file:

```
[mysqld]
myisam_sort_buffer_size=256M
key_buffer_size=1G
myisam_max_sort_file_size=100G
```

These settings do not include `read_buffer_size`. Setting `read_buffer_size` globally to a large value does so for all sessions and can cause performance to suffer due to excessive memory allocation for a server with many simultaneous sessions.

8.7 Optimizing for MEMORY Tables

Consider using `MEMORY` tables for noncritical data that is accessed often, and is read-only or rarely updated. Benchmark your application against equivalent `InnoDB` or `MyISAM` tables under a realistic workload, to confirm that any additional performance is worth the risk of losing data, or the overhead of copying data from a disk-based table at application start.

For best performance with `MEMORY` tables, examine the kinds of queries against each table, and specify the type to use for each associated index, either a B-tree index or a hash index. On the `CREATE INDEX` statement, use the clause `USING BTREE` or `USING HASH`. B-tree indexes are fast for queries that do greater-than or less-than comparisons through operators such as `>` or `BETWEEN`. Hash indexes are only fast for queries that look up single values through the `=` operator, or a restricted set of values through the `IN` operator. For why `USING BTREE` is often a better choice than the default `USING HASH`, see Section 8.2.1.19, “Avoiding Full Table Scans”. For implementation details of the different types of `MEMORY` indexes, see Section 8.3.8, “Comparison of B-Tree and Hash Indexes”.

8.8 Understanding the Query Execution Plan

Depending on the details of your tables, columns, indexes, and the conditions in your `WHERE` clause, the MySQL optimizer considers many techniques to efficiently perform the lookups involved in an SQL query. A query on a huge table can be performed without reading all the rows; a join involving several tables can be performed without comparing every combination of rows. The set of operations that the optimizer chooses to perform the most efficient query is called the “query execution plan”, also known as the `EXPLAIN` plan. Your goals are to recognize the aspects of the `EXPLAIN` plan that indicate a query is optimized well, and to learn the SQL syntax and indexing techniques to improve the plan if you see some inefficient operations.

8.8.1 Optimizing Queries with EXPLAIN

The `EXPLAIN` statement provides information about how MySQL executes statements:

- `EXPLAIN` works with `SELECT`, `DELETE`, `INSERT`, `REPLACE`, and `UPDATE` statements.
- When `EXPLAIN` is used with an explainable statement, MySQL displays information from the optimizer about the statement execution plan. That is, MySQL explains how it would process the statement, including information about how tables are joined and in which order. For information about using `EXPLAIN` to obtain execution plan information, see Section 8.8.2, “EXPLAIN Output Format”.
- When `EXPLAIN` is used with `FOR CONNECTION` *`connection_id`* rather than an explainable statement, it displays the execution plan for the statement executing in the named connection. See Section 8.8.4, “Obtaining Execution Plan Information for a Named Connection”.
- For `SELECT` statements, `EXPLAIN` produces additional execution plan information. See Section 8.8.3, “Extended EXPLAIN Output Format”.
- `EXPLAIN` is useful for examining queries involving partitioned tables. See Section 22.3.5, “Obtaining Information About Partitions”.

- The FORMAT option can be used to select the output format. TRADITIONAL presents the output in tabular format. This is the default if no FORMAT option is present. JSON format displays the information in JSON format.

With the help of EXPLAIN, you can see where you should add indexes to tables so that the statement executes faster by using indexes to find rows. You can also use EXPLAIN to check whether the optimizer joins the tables in an optimal order. To give a hint to the optimizer to use a join order corresponding to the order in which the tables are named in a SELECT statement, begin the statement with SELECT STRAIGHT_JOIN rather than just SELECT. (See Section 13.2.10, “SELECT Syntax”.) However, STRAIGHT_JOIN may prevent indexes from being used because it disables semi-join transformations. See Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”.

The optimizer trace may sometimes provide information complementary to that of EXPLAIN. However, the optimizer trace format and content are subject to change between versions. For details, see MySQL Internals: Tracing the Optimizer.

If you have a problem with indexes not being used when you believe that they should be, run ANALYZE TABLE to update table statistics, such as cardinality of keys, that can affect the choices the optimizer makes. See Section 13.7.2.1, “ANALYZE TABLE Syntax”.

Note

EXPLAIN can also be used to obtain information about the columns in a table. EXPLAIN *tbl_name* is synonymous with DESCRIBE *tbl_name* and SHOW COLUMNS FROM *tbl_name*. For more information, see Section 13.8.1, “DESCRIBE Syntax”, and Section 13.7.5.5, “SHOW COLUMNS Syntax”.

8.8.2 EXPLAIN Output Format

The EXPLAIN statement provides information about how MySQL executes statements. EXPLAIN works with SELECT, DELETE, INSERT, REPLACE, and UPDATE statements.

EXPLAIN returns a row of information for each table used in the SELECT statement. It lists the tables in the output in the order that MySQL would read them while processing the statement. MySQL resolves all joins using a nested-loop join method. This means that MySQL reads a row from the first table, and then finds a matching row in the second table, the third table, and so on. When all tables are processed, MySQL outputs the selected columns and backtracks through the table list until a table is found for which there are more matching rows. The next row is read from this table and the process continues with the next table.

Note

In older MySQL releases, partition and extended information was produced using EXPLAIN PARTITIONS and EXPLAIN EXTENDED. Those syntaxes are still recognized for backward compatibility but partition and extended output is now enabled by default, so the PARTITIONS and EXTENDED keywords are superfluous and deprecated. Their use results in a warning, and they will be removed from EXPLAIN syntax in a future MySQL release.

You cannot use the deprecated PARTITIONS and EXTENDED keywords together in the same EXPLAIN statement. In addition, neither of these keywords can be used together with the FORMAT option.

Note

MySQL Workbench has a Visual Explain capability that provides a visual representation of EXPLAIN output. See Tutorial: Using Visual Explain to Improve Query Performance.

- EXPLAIN Output Columns
- EXPLAIN Join Types
- EXPLAIN Extra Information
- EXPLAIN Output Interpretation

EXPLAIN Output Columns

This section describes the output columns produced by `EXPLAIN`. Later sections provide additional information about the `type` and `Extra` columns.

Each output row from `EXPLAIN` provides information about one table. Each row contains the values summarized in Table 8.1, “EXPLAIN Output Columns”, and described in more detail following the table. Column names are shown in the table's first column; the second column provides the equivalent property name shown in the output when `FORMAT=JSON` is used.

Table 8.1 EXPLAIN Output Columns

Column	JSON Name	Meaning
`id`	`select_id`	The `SELECT` identifier
`select_type`	None	The `SELECT` type
`table`	`table_name`	The table for the output row
`partitions`	`partitions`	The matching partitions
`type`	`access_type`	The join type
`possible_keys`	`possible_keys`	The possible indexes to choose
`key`	`key`	The index actually chosen
`key_len`	`key_length`	The length of the chosen key
`ref`	`ref`	The columns compared to the index
`rows`	`rows`	Estimate of rows to be examined
`filtered`	`filtered`	Percentage of rows filtered by table condition
`Extra`	None	Additional information

Note

JSON properties which are `NULL` are not displayed in JSON-formatted `EXPLAIN` output.

- `id` (JSON name: `select_id`)

  The `SELECT` identifier. This is the sequential number of the `SELECT` within the query. The value can be `NULL` if the row refers to the union result of other rows. In this case, the `table` column shows a value like `<unionM,N>` to indicate that the row refers to the union of the rows with `id` values of *`M`* and *`N`*.

- `select_type` (JSON name: none)

  The type of `SELECT`, which can be any of those shown in the following table. A JSON-formatted `EXPLAIN` exposes the `SELECT` type as a property of a `query_block`, unless it is `SIMPLE` or `PRIMARY`. The JSON names (where applicable) are also shown in the table.

`select_type` Value	JSON Name	Meaning
`SIMPLE`	None	Simple `SELECT` (not using `UNION` or subqueries)
`PRIMARY`	None	Outermost `SELECT`

`select_type` Value	JSON Name	Meaning
`UNION`	None	Second or later `SELECT` statement in a `UNION`
`DEPENDENT UNION`	`dependent` (`true`)	Second or later `SELECT` statement in a `UNION`, dependent on outer query
`UNION RESULT`	`union_result`	Result of a `UNION`.
`SUBQUERY`	None	First `SELECT` in subquery
`DEPENDENT SUBQUERY`	`dependent` (`true`)	First `SELECT` in subquery, dependent on outer query
`DERIVED`	None	Derived table `SELECT` (subquery in `FROM` clause)
`MATERIALIZED`	`materialized_from_subquery`	Materialized subquery
`UNCACHEABLE SUBQUERY`	`cacheable` (`false`)	A subquery for which the result cannot be cached and must be re-evaluated for each row of the outer query
`UNCACHEABLE UNION`	`cacheable` (`false`)	The second or later select in a `UNION` that belongs to an uncacheable subquery (see `UNCACHEABLE SUBQUERY`)

`DEPENDENT` typically signifies the use of a correlated subquery. See Section 13.2.11.7, “Correlated Subqueries”.

`DEPENDENT SUBQUERY` evaluation differs from `UNCACHEABLE SUBQUERY` evaluation. For `DEPENDENT SUBQUERY`, the subquery is re-evaluated only once for each set of different values of the variables from its outer context. For `UNCACHEABLE SUBQUERY`, the subquery is re-evaluated for each row of the outer context.

Cacheability of subqueries differs from caching of query results in the query cache (which is described in Section 8.10.3.1, “How the Query Cache Operates”). Subquery caching occurs during query execution, whereas the query cache is used to store results only after query execution finishes.

When you specify `FORMAT=JSON` with `EXPLAIN`, the output has no single property directly equivalent to `select_type`; the `query_block` property corresponds to a given `SELECT`. Properties equivalent to most of the `SELECT` subquery types just shown are available (an example being `materialized_from_subquery` for `MATERIALIZED`), and are displayed when appropriate. There are no JSON equivalents for `SIMPLE` or `PRIMARY`.

The `select_type` value for non-`SELECT` statements displays the statement type for affected tables. For example, `select_type` is `DELETE` for `DELETE` statements.

- `table` (JSON name: `table_name`)

  The name of the table to which the row of output refers. This can also be one of the following values:

  - `<union`*`M`*`,`*`N`*`>`: The row refers to the union of the rows with `id` values of *`M`* and *`N`*.
  - `<derived`*`N`*`>`: The row refers to the derived table result for the row with an `id` value of *`N`*. A derived table may result, for example, from a subquery in the `FROM` clause.
  - `<subquery`*`N`*`>`: The row refers to the result of a materialized subquery for the row with an `id` value of *`N`*. See Section 8.2.2.2, “Optimizing Subqueries with Materialization”.

- `partitions` (JSON name: `partitions`)

  The partitions from which records would be matched by the query. The value is `NULL` for nonpartitioned tables. See Section 22.3.5, “Obtaining Information About Partitions”.

- `type` (JSON name: `access_type`)

The join type. For descriptions of the different types, see `EXPLAIN` Join Types.

- `possible_keys` (JSON name: `possible_keys`)

  The `possible_keys` column indicates which indexes MySQL can choose from to find the rows in this table. Note that this column is totally independent of the order of the tables as displayed in the output from `EXPLAIN`. That means that some of the keys in `possible_keys` might not be usable in practice with the generated table order.

  If this column is `NULL` (or undefined in JSON-formatted output), there are no relevant indexes. In this case, you may be able to improve the performance of your query by examining the `WHERE` clause to check whether it refers to some column or columns that would be suitable for indexing. If so, create an appropriate index and check the query with `EXPLAIN` again. See Section 13.1.7, “ALTER TABLE Syntax”.

  To see what indexes a table has, use `SHOW INDEX FROM` *`tbl_name`*.

- `key` (JSON name: `key`)

  The `key` column indicates the key (index) that MySQL actually decided to use. If MySQL decides to use one of the `possible_keys` indexes to look up rows, that index is listed as the key value.

  It is possible that `key` will name an index that is not present in the `possible_keys` value. This can happen if none of the `possible_keys` indexes are suitable for looking up rows, but all the columns selected by the query are columns of some other index. That is, the named index covers the selected columns, so although it is not used to determine which rows to retrieve, an index scan is more efficient than a data row scan.

  For `InnoDB`, a secondary index might cover the selected columns even if the query also selects the primary key because `InnoDB` stores the primary key value with each secondary index. If `key` is `NULL`, MySQL found no index to use for executing the query more efficiently.

  To force MySQL to use or ignore an index listed in the `possible_keys` column, use `FORCE INDEX`, `USE INDEX`, or `IGNORE INDEX` in your query. See Section 8.9.4, “Index Hints”.

  For `MyISAM` tables, running `ANALYZE TABLE` helps the optimizer choose better indexes. For `MyISAM` tables, `myisamchk --analyze` does the same. See Section 13.7.2.1, “ANALYZE TABLE Syntax”, and Section 7.6, “MyISAM Table Maintenance and Crash Recovery”.

- `key_len` (JSON name: `key_length`)

  The `key_len` column indicates the length of the key that MySQL decided to use. The value of `key_len` enables you to determine how many parts of a multiple-part key MySQL actually uses. If the `key` column says `NULL`, the `len_len` column also says `NULL`.

  Due to the key storage format, the key length is one greater for a column that can be `NULL` than for a `NOT NULL` column.

- `ref` (JSON name: `ref`)

  The `ref` column shows which columns or constants are compared to the index named in the `key` column to select rows from the table.

  If the value is `func`, the value used is the result of some function. To see which function, use `SHOW WARNINGS` following `EXPLAIN` to see the extended `EXPLAIN` output. The function might actually be an operator such as an arithmetic operator.

- `rows` (JSON name: `rows`)

  The `rows` column indicates the number of rows MySQL believes it must examine to execute the query.

For `InnoDB` tables, this number is an estimate, and may not always be exact.

- `filtered` (JSON name: `filtered`)

  The `filtered` column indicates an estimated percentage of table rows that will be filtered by the table condition. That is, `rows` shows the estimated number of rows examined and `rows` × `filtered` / `100` shows the number of rows that will be joined with previous tables.

- `Extra` (JSON name: none)

  This column contains additional information about how MySQL resolves the query. For descriptions of the different values, see `EXPLAIN` Extra Information.

  There is no single JSON property corresponding to the `Extra` column; however, values that can occur in this column are exposed as JSON properties, or as the text of the `message` property.

EXPLAIN Join Types

The `type` column of `EXPLAIN` output describes how tables are joined. In JSON-formatted output, these are found as values of the `access_type` property. The following list describes the join types, ordered from the best type to the worst:

- `system`

  The table has only one row (= system table). This is a special case of the `const` join type.

- `const`

  The table has at most one matching row, which is read at the start of the query. Because there is only one row, values from the column in this row can be regarded as constants by the rest of the optimizer. `const` tables are very fast because they are read only once.

  `const` is used when you compare all parts of a `PRIMARY KEY` or `UNIQUE` index to constant values. In the following queries, *`tbl_name`* can be used as a `const` table:

```
SELECT * FROM tbl_name WHERE primary_key=1;

SELECT * FROM tbl_name
  WHERE primary_key_part1=1 AND primary_key_part2=2;
```

- `eq_ref`

  One row is read from this table for each combination of rows from the previous tables. Other than the `system` and `const` types, this is the best possible join type. It is used when all parts of an index are used by the join and the index is a `PRIMARY KEY` or `UNIQUE NOT NULL` index.

  `eq_ref` can be used for indexed columns that are compared using the `=` operator. The comparison value can be a constant or an expression that uses columns from tables that are read before this table. In the following examples, MySQL can use an `eq_ref` join to process *`ref_table`*:

```
SELECT * FROM ref_table,other_table
  WHERE ref_table.key_column=other_table.column;

SELECT * FROM ref_table,other_table
  WHERE ref_table.key_column_part1=other_table.column
  AND ref_table.key_column_part2=1;
```

- `ref`

  All rows with matching index values are read from this table for each combination of rows from the previous tables. `ref` is used if the join uses only a leftmost prefix of the key or if the key is not a

`PRIMARY KEY` or `UNIQUE` index (in other words, if the join cannot select a single row based on the key value). If the key that is used matches only a few rows, this is a good join type.

`ref` can be used for indexed columns that are compared using the `=` or `<=>` operator. In the following examples, MySQL can use a `ref` join to process *`ref_table`*:

```
SELECT * FROM ref_table WHERE key_column=expr;

SELECT * FROM ref_table,other_table
  WHERE ref_table.key_column=other_table.column;

SELECT * FROM ref_table,other_table
  WHERE ref_table.key_column_part1=other_table.column
  AND ref_table.key_column_part2=1;
```

- `fulltext`

  The join is performed using a `FULLTEXT` index.

- `ref_or_null`

  This join type is like `ref`, but with the addition that MySQL does an extra search for rows that contain `NULL` values. This join type optimization is used most often in resolving subqueries. In the following examples, MySQL can use a `ref_or_null` join to process *`ref_table`*:

  ```
  SELECT * FROM ref_table
    WHERE key_column=expr OR key_column IS NULL;
  ```

  See Section 8.2.1.12, “IS NULL Optimization”.

- `index_merge`

  This join type indicates that the Index Merge optimization is used. In this case, the `key` column in the output row contains a list of indexes used, and `key_len` contains a list of the longest key parts for the indexes used. For more information, see Section 8.2.1.3, “Index Merge Optimization”.

- `unique_subquery`

  This type replaces `eq_ref` for some `IN` subqueries of the following form:

  ```
  value IN (SELECT primary_key FROM single_table WHERE some_expr)
  ```

  `unique_subquery` is just an index lookup function that replaces the subquery completely for better efficiency.

- `index_subquery`

  This join type is similar to `unique_subquery`. It replaces `IN` subqueries, but it works for nonunique indexes in subqueries of the following form:

  ```
  value IN (SELECT key_column FROM single_table WHERE some_expr)
  ```

- `range`

  Only rows that are in a given range are retrieved, using an index to select the rows. The `key` column in the output row indicates which index is used. The `key_len` contains the longest key part that was used. The `ref` column is `NULL` for this type.

  `range` can be used when a key column is compared to a constant using any of the `=`, `<>`, `>`, `>=`, `<`, `<=`, `IS NULL`, `<=>`, `BETWEEN`, or `IN()` operators:

  ```
  ```

```
SELECT * FROM tbl_name
  WHERE key_column = 10;

SELECT * FROM tbl_name
  WHERE key_column BETWEEN 10 and 20;

SELECT * FROM tbl_name
  WHERE key_column IN (10,20,30);

SELECT * FROM tbl_name
  WHERE key_part1 = 10 AND key_part2 IN (10,20,30);
```

- `index`

  The `index` join type is the same as `ALL`, except that the index tree is scanned. This occurs two ways:

  - If the index is a covering index for the queries and can be used to satisfy all data required from the table, only the index tree is scanned. In this case, the `Extra` column says `Using index`. An index-only scan usually is faster than `ALL` because the size of the index usually is smaller than the table data.
  - A full table scan is performed using reads from the index to look up data rows in index order. `Uses index` does not appear in the `Extra` column.

  MySQL can use this join type when the query uses only columns that are part of a single index.

- `ALL`

  A full table scan is done for each combination of rows from the previous tables. This is normally not good if the table is the first table not marked `const`, and usually *very* bad in all other cases. Normally, you can avoid `ALL` by adding indexes that enable row retrieval from the table based on constant values or column values from earlier tables.

EXPLAIN Extra Information

The `Extra` column of `EXPLAIN` output contains additional information about how MySQL resolves the query. The following list explains the values that can appear in this column. Each item also indicates for JSON-formatted output which property displays the `Extra` value. For some of these, there is a specific property. The others display as the text of the `message` property.

If you want to make your queries as fast as possible, look out for `Extra` column values of `Using filesort` and `Using temporary`, or, in JSON-formatted `EXPLAIN` output, for `using_filesort` and `using_temporary_table` properties equal to `true`.

- `const row not found` (JSON property: `const_row_not_found`)

  For a query such as `SELECT ... FROM tbl_name`, the table was empty.

- `Deleting all rows` (JSON property: `message`)

  For `DELETE`, some storage engines (such as `MyISAM`) support a handler method that removes all table rows in a simple and fast way. This `Extra` value is displayed if the engine uses this optimization.

- `Distinct` (JSON property: `distinct`)

  MySQL is looking for distinct values, so it stops searching for more rows for the current row combination after it has found the first matching row.

- `FirstMatch(tbl_name)` (JSON property: `first_match`)

  The semi-join FirstMatch join shortcutting strategy is used for `tbl_name`.

- `Full scan on NULL key` (JSON property: `message`)

  This occurs for subquery optimization as a fallback strategy when the optimizer cannot use an index-lookup access method.

- `Impossible HAVING` (JSON property: `message`)

  The `HAVING` clause is always false and cannot select any rows.

- `Impossible WHERE` (JSON property: `message`)

  The `WHERE` clause is always false and cannot select any rows.

- `Impossible WHERE noticed after reading const tables` (JSON property: `message`)

  MySQL has read all `const` (and `system`) tables and notice that the `WHERE` clause is always false.

- `LooseScan(`*`m`*`..`*`n`*`)` (JSON property: `message`)

  The semi-join LooseScan strategy is used. *`m`* and *`n`* are key part numbers.

- `No matching min/max row` (JSON property: `message`)

  No row satisfies the condition for a query such as `SELECT MIN(...) FROM ... WHERE` *`condition`*.

- `no matching row in const table` (JSON property: `message`)

  For a query with a join, there was an empty table or a table with no rows satisfying a unique index condition.

- `No matching rows after partition pruning` (JSON property: `message`)

  For `DELETE` or `UPDATE`, the optimizer found nothing to delete or update after partition pruning. It is similar in meaning to `Impossible WHERE` for `SELECT` statements.

- `No tables used` (JSON property: `message`)

  The query has no `FROM` clause, or has a `FROM DUAL` clause.

  For `INSERT` or `REPLACE` statements, `EXPLAIN` displays this value when there is no `SELECT` part. For example, it appears for `EXPLAIN INSERT INTO t VALUES(10)` because that is equivalent to `EXPLAIN INSERT INTO t SELECT 10 FROM DUAL`.

- `Not exists` (JSON property: `message`)

  MySQL was able to do a `LEFT JOIN` optimization on the query and does not examine more rows in this table for the previous row combination after it finds one row that matches the `LEFT JOIN` criteria. Here is an example of the type of query that can be optimized this way:

  ```
  SELECT * FROM t1 LEFT JOIN t2 ON t1.id=t2.id
    WHERE t2.id IS NULL;
  ```

  Assume that `t2.id` is defined as `NOT NULL`. In this case, MySQL scans `t1` and looks up the rows in `t2` using the values of `t1.id`. If MySQL finds a matching row in `t2`, it knows that `t2.id` can never be `NULL`, and does not scan through the rest of the rows in `t2` that have the same `id` value. In other words, for each row in `t1`, MySQL needs to do only a single lookup in `t2`, regardless of how many rows actually match in `t2`.

- `Plan isn't ready yet` (JSON property: none)

  This value occurs with `EXPLAIN FOR CONNECTION` when the optimizer has not finished creating the execution plan for the statement executing in the named connection. If execution plan output

comprises multiple lines, any or all of them could have this `Extra` value, depending on the progress of the optimizer in determining the full execution plan.

- `Range checked for each record (index map: N)` (JSON property: `message`)

  MySQL found no good index to use, but found that some of indexes might be used after column values from preceding tables are known. For each row combination in the preceding tables, MySQL checks whether it is possible to use a `range` or `index_merge` access method to retrieve rows. This is not very fast, but is faster than performing a join with no index at all. The applicability criteria are as described in Section 8.2.1.2, “Range Optimization”, and Section 8.2.1.3, “Index Merge Optimization”, with the exception that all column values for the preceding table are known and considered to be constants.

  Indexes are numbered beginning with 1, in the same order as shown by `SHOW INDEX` for the table. The index map value *`N`* is a bitmask value that indicates which indexes are candidates. For example, a value of `0x19` (binary 11001) means that indexes 1, 4, and 5 will be considered.

- `Recursive` (JSON property: `recursive`)

  This indicates that the row applies to the recursive `SELECT` part of a recursive common table expression. See Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”.

- `Scanned N databases` (JSON property: `message`)

  This indicates how many directory scans the server performs when processing a query for `INFORMATION_SCHEMA` tables, as described in Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”. The value of *`N`* can be 0, 1, or `all`.

- `Select tables optimized away` (JSON property: `message`)

  The optimizer determined 1) that at most one row should be returned, and 2) that to produce this row, a deterministic set of rows must be read. When the rows to be read can be read during the optimization phase (for example, by reading index rows), there is no need to read any tables during query execution.

  The first condition is fulfilled when the query is implicitly grouped (contains an aggregate function but no `GROUP BY` clause). The second condition is fulfilled when one row lookup is performed per index used. The number of indexes read determines the number of rows to read.

  Consider the following implicitly grouped query:

  ```
  SELECT MIN(c1), MIN(c2) FROM t1;
  ```

  Suppose that `MIN(c1)` can be retrieved by reading one index row and `MIN(c2)` can be retrieved by reading one row from a different index. That is, for each column `c1` and `c2`, there exists an index where the column is the first column of the index. In this case, one row is returned, produced by reading two deterministic rows.

  This `Extra` value does not occur if the rows to read are not deterministic. Consider this query:

  ```
  SELECT MIN(c2) FROM t1 WHERE c1 <= 10;
  ```

  Suppose that `(c1, c2)` is a covering index. Using this index, all rows with `c1 <= 10` must be scanned to find the minimum `c2` value. By contrast, consider this query:

  ```
  SELECT MIN(c2) FROM t1 WHERE c1 = 10;
  ```

  In this case, the first index row with `c1 = 10` contains the minimum `c2` value. Only one row must be read to produce the returned row.

For storage engines that maintain an exact row count per table (such as `MyISAM`, but not `InnoDB`), this `Extra` value can occur for `COUNT(*)` queries for which the `WHERE` clause is missing or always true and there is no `GROUP BY` clause. (This is an instance of an implicitly grouped query where the storage engine influences whether a deterministic number of rows can be read.)

- `Skip_open_table`, `Open_frm_only`, `Open_full_table` (JSON property: `message`)

  These values indicate file-opening optimizations that apply to queries for `INFORMATION_SCHEMA` tables.

  - `Skip_open_table`: Table files do not need to be opened. The information is already available from the data dictionary.
  - `Open_frm_only`: Only the data dictionary need be read for table information.
  - `Open_full_table`: Unoptimized information lookup. Table information must be read from the data dictionary and by reading table files.

- `Start temporary`, `End temporary` (JSON property: `message`)

  This indicates temporary table use for the semi-join Duplicate Weedout strategy.

- `unique row not found` (JSON property: `message`)

  For a query such as `SELECT ... FROM tbl_name`, no rows satisfy the condition for a `UNIQUE` index or `PRIMARY KEY` on the table.

- `Using filesort` (JSON property: `using_filesort`)

  MySQL must do an extra pass to find out how to retrieve the rows in sorted order. The sort is done by going through all rows according to the join type and storing the sort key and pointer to the row for all rows that match the `WHERE` clause. The keys then are sorted and the rows are retrieved in sorted order. See Section 8.2.1.13, “ORDER BY Optimization”.

- `Using index` (JSON property: `using_index`)

  The column information is retrieved from the table using only information in the index tree without having to do an additional seek to read the actual row. This strategy can be used when the query uses only columns that are part of a single index.

  For `InnoDB` tables that have a user-defined clustered index, that index can be used even when `Using index` is absent from the `Extra` column. This is the case if `type` is `index` and `key` is `PRIMARY`.

- `Using index condition` (JSON property: `using_index_condition`)

  Tables are read by accessing index tuples and testing them first to determine whether to read full table rows. In this way, index information is used to defer (“push down”) reading full table rows unless it is necessary. See Section 8.2.1.5, “Index Condition Pushdown Optimization”.

- `Using index for group-by` (JSON property: `using_index_for_group_by`)

  Similar to the `Using index` table access method, `Using index for group-by` indicates that MySQL found an index that can be used to retrieve all columns of a `GROUP BY` or `DISTINCT` query without any extra disk access to the actual table. Additionally, the index is used in the most efficient way so that for each group, only a few index entries are read. For details, see Section 8.2.1.14, “GROUP BY Optimization”.

- `Using join buffer (Block Nested Loop)`, `Using join buffer (Batched Key Access)` (JSON property: `using_join_buffer`)

Tables from earlier joins are read in portions into the join buffer, and then their rows are used from the buffer to perform the join with the current table. `(Block Nested Loop)` indicates use of the Block Nested-Loop algorithm and `(Batched Key Access)` indicates use of the Batched Key Access algorithm. That is, the keys from the table on the preceding line of the `EXPLAIN` output will be buffered, and the matching rows will be fetched in batches from the table represented by the line in which `Using join buffer` appears.

In JSON-formatted output, the value of `using_join_buffer` is always either one of `Block Nested Loop` or `Batched Key Access`.

- `Using MRR` (JSON property: `message`)

  Tables are read using the Multi-Range Read optimization strategy. See Section 8.2.1.10, “Multi-Range Read Optimization”.

- `Using sort_union(...)`, `Using union(...)`, `Using intersect(...)` (JSON property: `message`)

  These indicate the particular algorithm showing how index scans are merged for the `index_merge` join type. See Section 8.2.1.3, “Index Merge Optimization”.

- `Using temporary` (JSON property: `using_temporary_table`)

  To resolve the query, MySQL needs to create a temporary table to hold the result. This typically happens if the query contains `GROUP BY` and `ORDER BY` clauses that list columns differently.

- `Using where` (JSON property: `attached_condition`)

  A `WHERE` clause is used to restrict which rows to match against the next table or send to the client. Unless you specifically intend to fetch or examine all rows from the table, you may have something wrong in your query if the `Extra` value is not `Using where` and the table join type is `ALL` or `index`.

  `Using where` has no direct counterpart in JSON-formatted output; the `attached_condition` property contains any `WHERE` condition used.

- `Using where with pushed condition` (JSON property: `message`)

  This item applies to `NDB` tables *only*. It means that MySQL Cluster is using the Condition Pushdown optimization to improve the efficiency of a direct comparison between a nonindexed column and a constant. In such cases, the condition is “pushed down” to the cluster's data nodes and is evaluated on all data nodes simultaneously. This eliminates the need to send nonmatching rows over the network, and can speed up such queries by a factor of 5 to 10 times over cases where Condition Pushdown could be but is not used. For more information, see Section 8.2.1.4, “Engine Condition Pushdown Optimization”.

- `Zero limit` (JSON property: `message`)

  The query had a `LIMIT 0` clause and cannot select any rows.

EXPLAIN Output Interpretation

You can get a good indication of how good a join is by taking the product of the values in the `rows` column of the `EXPLAIN` output. This should tell you roughly how many rows MySQL must examine to execute the query. If you restrict queries with the `max_join_size` system variable, this row product also is used to determine which multiple-table `SELECT` statements to execute and which to abort. See Section 5.1.1, “Configuring the Server”.

The following example shows how a multiple-table join can be optimized progressively based on the information provided by `EXPLAIN`.

Suppose that you have the SELECT statement shown here and that you plan to examine it using EXPLAIN:

```
EXPLAIN SELECT tt.TicketNumber, tt.TimeIn,
               tt.ProjectReference, tt.EstimatedShipDate,
               tt.ActualShipDate, tt.ClientID,
               tt.ServiceCodes, tt.RepetitiveID,
               tt.CurrentProcess, tt.CurrentDPPerson,
               tt.RecordVolume, tt.DPPrinted, et.COUNTRY,
               et_1.COUNTRY, do.CUSTNAME
        FROM tt, et, et AS et_1, do
        WHERE tt.SubmitTime IS NULL
          AND tt.ActualPC = et.EMPLOYID
          AND tt.AssignedPC = et_1.EMPLOYID
          AND tt.ClientID = do.CUSTNMBR;
```

For this example, make the following assumptions:

- The columns being compared have been declared as follows.

Table	Column	Data Type
tt	ActualPC	CHAR(10)
tt	AssignedPC	CHAR(10)
tt	ClientID	CHAR(10)
et	EMPLOYID	CHAR(15)
do	CUSTNMBR	CHAR(15)

- The tables have the following indexes.

Table	Index
tt	ActualPC
tt	AssignedPC
tt	ClientID
et	EMPLOYID (primary key)
do	CUSTNMBR (primary key)

- The tt.ActualPC values are not evenly distributed.

Initially, before any optimizations have been performed, the EXPLAIN statement produces the following information:

```
table type possible_keys key  key_len ref  rows  Extra
et    ALL  PRIMARY       NULL NULL    NULL 74
do    ALL  PRIMARY       NULL NULL    NULL 2135
et_1  ALL  PRIMARY       NULL NULL    NULL 74
tt    ALL  AssignedPC,   NULL NULL    NULL 3872
           ClientID,
           ActualPC
      Range checked for each record (index map: 0x23)
```

Because type is ALL for each table, this output indicates that MySQL is generating a Cartesian product of all the tables; that is, every combination of rows. This takes quite a long time, because the product of the number of rows in each table must be examined. For the case at hand, this product is 74 × 2135 × 74 × 3872 = 45,268,558,720 rows. If the tables were bigger, you can only imagine how long it would take.

One problem here is that MySQL can use indexes on columns more efficiently if they are declared as the same type and size. In this context, VARCHAR and CHAR are considered the same if they are

declared as the same size. `tt.ActualPC` is declared as `CHAR(10)` and `et.EMPLOYID` is `CHAR(15)`, so there is a length mismatch.

To fix this disparity between column lengths, use `ALTER TABLE` to lengthen `ActualPC` from 10 characters to 15 characters:

```
mysql> ALTER TABLE tt MODIFY ActualPC VARCHAR(15);
```

Now `tt.ActualPC` and `et.EMPLOYID` are both `VARCHAR(15)`. Executing the `EXPLAIN` statement again produces this result:

```
table type   possible_keys key     key_len ref         rows    Extra
tt    ALL    AssignedPC,   NULL    NULL    NULL        3872    Using
             ClientID,                                        where
             ActualPC
do    ALL    PRIMARY       NULL    NULL    NULL        2135
      Range checked for each record (index map: 0x1)
et_1  ALL    PRIMARY       NULL    NULL    NULL        74
      Range checked for each record (index map: 0x1)
et    eq_ref PRIMARY       PRIMARY 15      tt.ActualPC 1
```

This is not perfect, but is much better: The product of the `rows` values is less by a factor of 74. This version executes in a couple of seconds.

A second alteration can be made to eliminate the column length mismatches for the `tt.AssignedPC = et_1.EMPLOYID` and `tt.ClientID = do.CUSTNMBR` comparisons:

```
mysql> ALTER TABLE tt MODIFY AssignedPC VARCHAR(15),
                      MODIFY ClientID   VARCHAR(15);
```

After that modification, `EXPLAIN` produces the output shown here:

```
table type   possible_keys key      key_len ref           rows Extra
et    ALL    PRIMARY       NULL     NULL    NULL          74
tt    ref    AssignedPC,   ActualPC 15      et.EMPLOYID   52   Using
             ClientID,                                        where
             ActualPC
et_1  eq_ref PRIMARY       PRIMARY  15      tt.AssignedPC 1
do    eq_ref PRIMARY       PRIMARY  15      tt.ClientID   1
```

At this point, the query is optimized almost as well as possible. The remaining problem is that, by default, MySQL assumes that values in the `tt.ActualPC` column are evenly distributed, and that is not the case for the `tt` table. Fortunately, it is easy to tell MySQL to analyze the key distribution:

```
mysql> ANALYZE TABLE tt;
```

With the additional index information, the join is perfect and `EXPLAIN` produces this result:

```
table type   possible_keys key     key_len ref           rows Extra
tt    ALL    AssignedPC    NULL    NULL    NULL          3872 Using
             ClientID,                                       where
             ActualPC
et    eq_ref PRIMARY       PRIMARY 15      tt.ActualPC   1
et_1  eq_ref PRIMARY       PRIMARY 15      tt.AssignedPC 1
do    eq_ref PRIMARY       PRIMARY 15      tt.ClientID   1
```

The `rows` column in the output from `EXPLAIN` is an educated guess from the MySQL join optimizer. Check whether the numbers are even close to the truth by comparing the `rows` product with the actual number of rows that the query returns. If the numbers are quite different, you might get better performance by using `STRAIGHT_JOIN` in your `SELECT` statement and trying to list the tables in a different order in the `FROM` clause. (However, `STRAIGHT_JOIN` may prevent indexes from being used

because it disables semi-join transformations. See Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”.)

It is possible in some cases to execute statements that modify data when `EXPLAIN SELECT` is used with a subquery; for more information, see Section 13.2.11.8, “Derived Tables (Subqueries in the FROM Clause)”.

8.8.3 Extended EXPLAIN Output Format

For `SELECT` statements, the `EXPLAIN` statement produces extra (“extended”) information that is not part of `EXPLAIN` output but can be viewed by issuing a `SHOW WARNINGS` statement following `EXPLAIN`. The `Message` value in `SHOW WARNINGS` output displays how the optimizer qualifies table and column names in the `SELECT` statement, what the `SELECT` looks like after the application of rewriting and optimization rules, and possibly other notes about the optimization process.

The extended information displayable with a `SHOW WARNINGS` statement following `EXPLAIN` is produced only for `SELECT` statements. `SHOW WARNINGS` displays an empty result for other explainable statements (`DELETE`, `INSERT`, `REPLACE`, and `UPDATE`).

Note

In older MySQL releases, extended information was produced using `EXPLAIN EXTENDED`. That syntax is still recognized for backward compatibility but extended output is now enabled by default, so the `EXTENDED` keyword is superfluous and deprecated. Its use results in a warning, and it will be removed from `EXPLAIN` syntax in a future MySQL release.

Here is an example of extended `EXPLAIN` output:

```
mysql> EXPLAIN
       SELECT t1.a, t1.a IN (SELECT t2.a FROM t2) FROM t1\G
*************************** 1. row ***************************
           id: 1
  select_type: PRIMARY
        table: t1
         type: index
possible_keys: NULL
          key: PRIMARY
      key_len: 4
          ref: NULL
         rows: 4
     filtered: 100.00
        Extra: Using index
*************************** 2. row ***************************
           id: 2
  select_type: SUBQUERY
        table: t2
         type: index
possible_keys: a
          key: a
      key_len: 5
          ref: NULL
         rows: 3
     filtered: 100.00
        Extra: Using index
2 rows in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: /* select#1 */ select `test`.`t1`.`a` AS `a`,
         <in_optimizer>(`test`.`t1`.`a`,`test`.`t1`.`a` in
         ( <materialize> (/* select#2 */ select `test`.`t2`.`a`
         from `test`.`t2` where 1 having 1 ),
         <primary_index_lookup>(`test`.`t1`.`a` in
```

```
          <temporary table> on <auto_key>
          where ((`test`.`t1`.`a` = `materialized-subquery`.`a`))))) AS `t1.a
          IN (SELECT t2.a FROM t2)` from `test`.`t1`
1 row in set (0.00 sec)
```

Because the statement displayed by `SHOW WARNINGS` may contain special markers to provide information about query rewriting or optimizer actions, the statement is not necessarily valid SQL and is not intended to be executed. The output may also include rows with `Message` values that provide additional non-SQL explanatory notes about actions taken by the optimizer.

The following list describes special markers that can appear in the extended output displayed by `SHOW WARNINGS`:

- `<auto_key>`

  An automatically generated key for a temporary table.

- `<cache>(expr)`

  The expression (such as a scalar subquery) is executed once and the resulting value is saved in memory for later use. For results consisting of multiple values, a temporary table may be created and you will see `<temporary table>` instead.

- `<exists>(query fragment)`

  The subquery predicate is converted to an `EXISTS` predicate and the subquery is transformed so that it can be used together with the `EXISTS` predicate.

- `<in_optimizer>(query fragment)`

  This is an internal optimizer object with no user significance.

- `<index_lookup>(query fragment)`

  The query fragment is processed using an index lookup to find qualifying rows.

- `<if>(condition, expr1, expr2)`

  If the condition is true, evaluate to *`expr1`*, otherwise *`expr2`*.

- `<is_not_null_test>(expr)`

  A test to verify that the expression does not evaluate to `NULL`.

- `<materialize>(query fragment)`

  Subquery materialization is used.

- `` `materialized-subquery`.col_name ``

  A reference to the column *`col_name`* in an internal temporary table materialized to hold the result from evaluating a subquery.

- `<primary_index_lookup>(query fragment)`

  The query fragment is processed using a primary key lookup to find qualifying rows.

- `<ref_null_helper>(expr)`

  This is an internal optimizer object with no user significance.

- `/* select#N */ select_stmt`

  The `SELECT` is associated with the row in non-extended `EXPLAIN` output that has an `id` value of *`N`*.

- *outer_tables* semi join (*inner_tables*)

  A semi-join operation. *inner_tables* shows the tables that were not pulled out. See Section 8.2.2.1, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions with Semi-Join Transformations”.

- <temporary table>

  This represents an internal temporary table created to cache an intermediate result.

When some tables are of const or system type, expressions involving columns from these tables are evaluated early by the optimizer and are not part of the displayed statement. However, with FORMAT=JSON, some const table accesses are displayed as a ref access that uses a const value.

8.8.4 Obtaining Execution Plan Information for a Named Connection

To obtain the execution plan for an explainable statement executing in a named connection, use this statement:

```
EXPLAIN [options] FOR CONNECTION connection_id;
```

EXPLAIN FOR CONNECTION returns the EXPLAIN information that is currently being used to execute a query in a given connection. Because of changes to data (and supporting statistics) it may produce a different result from running EXPLAIN on the equivalent query text. This difference in behavior can be useful in diagnosing more transient performance problems. For example, if you are running a statement in one session that is taking a long time to complete, using EXPLAIN FOR CONNECTION in another session may yield useful information about the cause of the delay.

connection_id is the connection identifier, as obtained from the INFORMATION_SCHEMA PROCESSLIST table or the SHOW PROCESSLIST statement. If you have the PROCESS privilege, you can specify the identifier for any connection. Otherwise, you can specify the identifier only for your own connections.

If the named connection is not executing a statement, the result is empty. Otherwise, EXPLAIN FOR CONNECTION applies only if the statement being executed in the named connection is explainable. This includes SELECT, DELETE, INSERT, REPLACE, and UPDATE. (However, EXPLAIN FOR CONNECTION does not work for prepared statements, even prepared statements of those types.)

If the named connection is executing an explainable statement, the output is what you would obtain by using EXPLAIN on the statement itself.

If the named connection is executing a statement that is not explainable, an error occurs. For example, you cannot name the connection identifier for your current session because EXPLAIN is not explainable:

```
mysql> SELECT CONNECTION_ID();
+-----------------+
| CONNECTION_ID() |
+-----------------+
|             373 |
+-----------------+
1 row in set (0.00 sec)

mysql> EXPLAIN FOR CONNECTION 373;
ERROR 1889 (HY000): EXPLAIN FOR CONNECTION command is supported
only for SELECT/UPDATE/INSERT/DELETE/REPLACE
```

The Com_explain_other status variable indicates the number of EXPLAIN FOR CONNECTION statements executed.

8.8.5 Estimating Query Performance

In most cases, you can estimate query performance by counting disk seeks. For small tables, you can usually find a row in one disk seek (because the index is probably cached). For bigger tables, you can estimate that, using B-tree indexes, you need this many seeks to find a row: `log(row_count) / log(index_block_length / 3 * 2 / (index_length + data_pointer_length)) + 1`.

In MySQL, an index block is usually 1,024 bytes and the data pointer is usually four bytes. For a 500,000-row table with a key value length of three bytes (the size of `MEDIUMINT`), the formula indicates `log(500,000)/log(1024/3*2/(3+4)) + 1` = `4` seeks.

This index would require storage of about 500,000 * 7 * 3/2 = 5.2MB (assuming a typical index buffer fill ratio of 2/3), so you probably have much of the index in memory and so need only one or two calls to read data to find the row.

For writes, however, you need four seek requests to find where to place a new index value and normally two seeks to update the index and write the row.

The preceding discussion does not mean that your application performance slowly degenerates by log *N*. As long as everything is cached by the OS or the MySQL server, things become only marginally slower as the table gets bigger. After the data gets too big to be cached, things start to go much slower until your applications are bound only by disk seeks (which increase by log *N*). To avoid this, increase the key cache size as the data grows. For `MyISAM` tables, the key cache size is controlled by the `key_buffer_size` system variable. See Section 5.1.1, “Configuring the Server”.

8.9 Controlling the Query Optimizer

MySQL provides optimizer control through system variables that affect how query plans are evaluated, switchable optimizations, optimizer and index hints, and the optimizer cost model.

The server also maintains statistics about column values, although the optimizer does not yet use this information.

8.9.1 Controlling Query Plan Evaluation

The task of the query optimizer is to find an optimal plan for executing an SQL query. Because the difference in performance between “good” and “bad” plans can be orders of magnitude (that is, seconds versus hours or even days), most query optimizers, including that of MySQL, perform a more or less exhaustive search for an optimal plan among all possible query evaluation plans. For join queries, the number of possible plans investigated by the MySQL optimizer grows exponentially with the number of tables referenced in a query. For small numbers of tables (typically less than 7 to 10) this is not a problem. However, when larger queries are submitted, the time spent in query optimization may easily become the major bottleneck in the server's performance.

A more flexible method for query optimization enables the user to control how exhaustive the optimizer is in its search for an optimal query evaluation plan. The general idea is that the fewer plans that are investigated by the optimizer, the less time it spends in compiling a query. On the other hand, because the optimizer skips some plans, it may miss finding an optimal plan.

The behavior of the optimizer with respect to the number of plans it evaluates can be controlled using two system variables:

- The `optimizer_prune_level` variable tells the optimizer to skip certain plans based on estimates of the number of rows accessed for each table. Our experience shows that this kind of “educated guess” rarely misses optimal plans, and may dramatically reduce query compilation times. That is why this option is on (`optimizer_prune_level=1`) by default. However, if you believe that the optimizer missed a better query plan, this option can be switched off (`optimizer_prune_level=0`) with the risk that query compilation may take much longer. Note that, even with the use of this heuristic, the optimizer still explores a roughly exponential number of plans.

- The `optimizer_search_depth` variable tells how far into the “future” of each incomplete plan the optimizer should look to evaluate whether it should be expanded further. Smaller values of `optimizer_search_depth` may result in orders of magnitude smaller query compilation times. For example, queries with 12, 13, or more tables may easily require hours and even days to compile if `optimizer_search_depth` is close to the number of tables in the query. At the same time, if compiled with `optimizer_search_depth` equal to 3 or 4, the optimizer may compile in less than a minute for the same query. If you are unsure of what a reasonable value is for `optimizer_search_depth`, this variable can be set to 0 to tell the optimizer to determine the value automatically.

8.9.2 Optimizer Hints

One means of control over optimizer strategies is to set the `optimizer_switch` system variable (see Section 8.9.3, “Switchable Optimizations”). Changes to this variable affect execution of all subsequent queries; to affect one query differently from another, it's necessary to change `optimizer_switch` before each one.

Another way to control the optimizer is by using optimizer hints, which can be specified within individual statements. Because optimizer hints apply on a per-statement basis, they provide finer control over statement execution plans than can be achieved using `optimizer_switch`. For example, you can enable an optimization for one table in a statement and disable the optimization for a different table. Hints within a statement take precedence over `optimizer_switch` flags.

Examples:

```
SELECT /*+ NO_RANGE_OPTIMIZATION(t3 PRIMARY, f2_idx) */ f1
  FROM t3 WHERE f1 > 30 AND f1 < 33;
SELECT /*+ BKA(t1) NO_BKA(t2) */ * FROM t1 INNER JOIN t2 WHERE ...;
SELECT /*+ NO_ICP(t1, t2) */ * FROM t1 INNER JOIN t2 WHERE ...;
SELECT /*+ SEMIJOIN(FIRSTMATCH, LOOSESCAN) */ * FROM t1 ...;
EXPLAIN SELECT /*+ NO_ICP(t1) */ * FROM t1 WHERE ...;
SELECT /*+ MERGE(dt) */ * FROM (SELECT * FROM t1) AS dt;
```

Optimizer hints, described here, differ from index hints, described in Section 8.9.4, “Index Hints”. Optimizer and index hints may be used separately or together.

- Optimizer Hint Overview
- Optimizer Hint Syntax
- Join-Order Optimizer Hints
- Table-Level Optimizer Hints
- Index-Level Optimizer Hints
- Subquery Optimizer Hints
- Statement Execution Time Optimizer Hints
- Optimizer Hints for Naming Query Blocks

Optimizer Hint Overview

Optimizer hints apply at different scope levels:

- Global: The hint affects the entire statement
- Query block: The hint affects a particular query block within a statement
- Table-level: The hint affects a particular table within a query block

- Index-level: The hint affects a particular index within a table

The following table summarizes the available optimizer hints, the optimizer strategies they affect, and the scope or scopes at which they apply. More details are given later.

Table 8.2 Optimizer Hints Available

Hint Name	Description	Applicable Scopes
BKA, NO_BKA	Affects Batched Key Access join processing	Query block, table
BNL, NO_BNL	Affects Block Nested-Loop join processing	Query block, table
INDEX_MERGE, NO_INDEX_MERGE	Affects Index Merge optimization	Table, index
JOIN_FIXED_ORDER	Use table order specified in FROM clause for join order	Query block
JOIN_ORDER	Use table order specified in hint for join order	Query block
JOIN_PREFIX	Use table order specified in hint for first tables of join order	Query block
JOIN_SUFFIX	Use table order specified in hint for last tables of join order	Query block
MAX_EXECUTION_TIME	Limits statement execution time	Global
MERGE, NO_MERGE	Affects derived table/view merging into outer query block	Table
MRR, NO_MRR	Affects Multi-Range Read optimization	Table, index
NO_ICP	Affects Index Condition Pushdown optimization	Table, index
NO_RANGE_OPTIMIZATION	Affects range optimization	Table, index
QB_NAME	Assigns name to query block	Query block
SEMIJOIN, NO_SEMIJOIN	Affects semi-join strategies	Query block
SUBQUERY	Affects materialization, IN-to-EXISTS subquery stratgies	Query block

Disabling an optimization prevents the optimizer from using it. Enabling an optimization means the optimizer is free to use the strategy if it applies to statement execution, not that the optimizer necessarily will use it.

Optimizer Hint Syntax

MySQL supports comments in SQL statements as described in Section 9.6, “Comment Syntax”. Optimizer hints must be specified within `/*+ ... */` comments. That is, optimizer hints use a variant of `/* ... */` C-style comment syntax, with a `+` character following the `/*` comment opening sequence. Examples:

```
/*+ BKA(t1) */
/*+ BNL(t1, t2) */
/*+ NO_RANGE_OPTIMIZATION(t4 PRIMARY) */
/*+ QB_NAME(qb2) */
```

Whitespace is permitted after the `+` character.

The parser recognizes optimizer hint comments after the initial keyword of `SELECT`, `UPDATE`, `INSERT`, `REPLACE`, and `DELETE` statements. Hints are permitted in these contexts:

- At the beginning of query and data change statements:

```
SELECT /*+ ... */ ...
INSERT /*+ ... */ ...
REPLACE /*+ ... */ ...
UPDATE /*+ ... */ ...
DELETE /*+ ... */ ...
```

- At the beginning of query blocks:

```
(SELECT /*+ ... */ ... )
(SELECT ... ) UNION (SELECT /*+ ... */ ... )
(SELECT /*+ ... */ ... ) UNION (SELECT /*+ ... */ ... )
UPDATE ... WHERE x IN (SELECT /*+ ... */ ...)
INSERT ... SELECT /*+ ... */ ...
```

- In hintable statements prefaced by `EXPLAIN`. For example:

```
EXPLAIN SELECT /*+ ... */ ...
EXPLAIN UPDATE ... WHERE x IN (SELECT /*+ ... */ ...)
```

  The implication is that you can use `EXPLAIN` to see how optimizer hints affect execution plans. Use `SHOW WARNINGS` immediately after `EXPLAIN` to see how hints are used. The extended `EXPLAIN` output displayed by a following `SHOW WARNINGS` indicates which hints were used. Ignored hints are not displayed.

A hint comment may contain multiple hints, but a query block cannot contain multiple hint comments. This is valid:

```
SELECT /*+ BNL(t1) BKA(t2) */ ...
```

But this is invalid:

```
SELECT /*+ BNL(t1) */ /* BKA(t2) */ ...
```

When a hint comment contains multiple hints, the possibility of duplicates and conflicts exists. The following general guidelines apply. For specific hint types, additional rules may apply, as indicated in the hint descriptions.

- Duplicate hints: For a hint such as `/*+ MRR(idx1) MRR(idx1) */`, MySQL uses the first hint and issues a warning about the duplicate hint.
- Conflicting hints: For a hint such as `/*+ MRR(idx1) NO_MRR(idx1) */`, MySQL uses the first hint and issues a warning about the second conflicting hint.

Query block names are identifiers and follow the usual rules about what names are valid and how to quote them (see Section 9.2, “Schema Object Names”).

Hint names, query block names, and strategy names are not case sensitive. References to table and index names follow the usual identifier case sensitivity rules (see Section 9.2.2, “Identifier Case Sensitivity”).

Join-Order Optimizer Hints

Join-order hints affect the order in which the optimizer joins tables.

Syntax of the `JOIN_FIXED_ORDER` hint:

```
hint_name([@query_block_name])
```

Syntax of other join-order hints:

```
hint_name([@query_block_name] tbl_name [, tbl_name] ...)
hint_name(tbl_name[@query_block_name] [, tbl_name[@query_block_name]] ...)
```

The syntax refers to these terms:

- *`hint_name`*: These hint names are permitted:
  - `JOIN_FIXED_ORDER`: Force the optimizer to join tables using the order in which they appear in the `FROM` clause. This is the same as specifying `SELECT STRAIGHT_JOIN`.
  - `JOIN_ORDER`: Instruct the optimizer to join tables using the specified table order. The hint applies to the named tables. The optimizer may place tables that are not named anywhere in the join order, including between specified tables.
  - `JOIN_PREFIX`: Instruct the optimizer to join tables using the specified table order for the first tables of the join execution plan. The hint applies to the named tables. The optimizer places all other tables after the named tables.
  - `JOIN_SUFFIX`: Instruct the optimizer to join tables using the specified table order for the last tables of the join execution plan. The hint applies to the named tables. The optimizer places all other tables before the named tables.
- *`tbl_name`*: The name of a table used in the statement. A hint that names tables applies to all tables that it names. The `JOIN_FIXED_ORDER` hint names no tables and applies to all tables in the `FROM` clause of the query block in which it occurs.

  If a table has an alias, hints must refer to the alias, not the table name.

  Table names in hints cannot be qualified with schema names.
- *`query_block_name`*: The query block to which the hint applies. If the hint includes no leading *`@query_block_name`*, the hint applies to the query block in which it occurs. For *`tbl_name@query_block_name`* syntax, the hint applies to the named table in the named query block. To assign a name to a query block, see Optimizer Hints for Naming Query Blocks.

Example:

```
SELECT
/*+ JOIN_PREFIX(t2, t5@subq2, t4@subq1)
    JOIN_ORDER(t4@subq1, t3)
    JOIN_SUFFIX(t1) */
COUNT(*) FROM t1 JOIN t2 JOIN t3
           WHERE t1.f1 IN (SELECT /*+ QB_NAME(subq1) */ f1 FROM t4)
             AND t2.f1 IN (SELECT /*+ QB_NAME(subq2) */ f1 FROM t5);
```

Hints control the behavior of semi-join tables that are merged to the outer query block. If subqueries `subq1` and `subq2` are converted to semi-joins, tables `t4@subq1` and `t5@subq2` are merged to the outer query block. In this case, the hint specified in the outer query block controls the behavior of `t4@subq1`, `t5@subq2` tables.

The optimizer resolves join-order hints according to these principles:

- Multiple hint instances

  Only one `JOIN_PREFIX` and `JOIN_SUFFIX` hint of each type are applied. Any later hints of the same type are ignored with a warning. `JOIN_ORDER` can be specified several times.

  Examples:

```
/*+ JOIN_PREFIX(t1) JOIN_PREFIX(t2) */
```

The second `JOIN_PREFIX` hint is ignored with a warning.

```
/*+ JOIN_PREFIX(t1) JOIN_SUFFIX(t2) */
```

Both hints are applicable. No warning occurs.

```
/*+ JOIN_ORDER(t1, t2) JOIN_ORDER(t2, t3) */
```

Both hints are applicable. No warning occurs.

- Conflicting hints

  In some cases hints can conflict, such as when `JOIN_ORDER` and `JOIN_PREFIX` have table orders that are impossible to apply at the same time:

  ```
  SELECT /*+ JOIN_ORDER(t1, t2) JOIN_PREFIX(t2, t1) */ ... FROM t1, t2;
  ```

  In this case, the first specified hint is applied and subsequent conflicting hints are ignored with no warning. A valid hint that is impossible to apply is silently ignored with no warning.

- Ignored hints

  A hint is ignored if a table specified in the hint has a circular dependency.

  Example:

  ```
  /*+ JOIN_ORDER(t1, t2) JOIN_PREFIX(t2, t1) */
  ```

  The `JOIN_ORDER` hint sets table `t2` dependent on `t1`. The `JOIN_PREFIX` hint is ignored because table `t1` cannot be dependent on `t2`. Ignored hints are not displayed in extended `EXPLAIN` output.

- Interaction with `const` tables

  The MySQL optimizer places `const` tables first in the join order, and the position of a `const` table cannot be affected by hints. References to `const` tables in join-order hints are ignored, although the hint is still applicable. For example, these are equivalent:

  ```
  JOIN_ORDER(t1, const_tbl, t2)
  JOIN_ORDER(t1, t2)
  ```

  Accepted hints shown in extended `EXPLAIN` output include `const` tables as they were specified.

- Interaction with types of join operations

  MySQL supports several type of joins: `LEFT`, `RIGHT`, `INNER`, `CROSS`, `STRAIGHT_JOIN`. A hint that conflicts with the specified type of join is ignored with no warning.

  Example:

  ```
  SELECT /*+ JOIN_PREFIX(t1, t2) */FROM t2 LEFT JOIN t1;
  ```

  Here a conflict occurs between the requested join order in the hint and the order required by the `LEFT JOIN`. The hint is ignored with no warning.

Table-Level Optimizer Hints

Table-level hints affect:

- Use of the Block Nested-Loop (BNL) and Batched Key Access (BKA) join-processing algorithms (see Section 8.2.1.11, “Block Nested-Loop and Batched Key Access Joins”).
- Whether derived tables, view references, or common table expressions should be merged into the outer query block, or materialized using an internal temporary table.

These hint types apply to specific tables, or all tables in a query block.

Syntax of table-level hints:

```
hint_name([@query_block_name] [tbl_name [, tbl_name] ...])
hint_name([tbl_name@query_block_name [, tbl_name@query_block_name] ...])
```

The syntax refers to these terms:

- *`hint_name`*: These hint names are permitted:
  - `BKA`, `NO_BKA`: Enable or disable BKA for the specified tables.
  - `BNL`, `NO_BNL`: Enable or disable BNL for the specified tables.
  - `MERGE`, `NO_MERGE`: Enable merging for the specified tables, view references or common table expressions; or disable merging and use materialization instead.

Note

To use a BNL or BKA hint to enable join buffering for any inner table of an outer join, join buffering must be enabled for all inner tables of the outer join.

- *`tbl_name`*: The name of a table used in the statement. The hint applies to all tables that it names. If the hint names no tables, it applies to all tables of the query block in which it occurs.

  If a table has an alias, hints must refer to the alias, not the table name.

  Table names in hints cannot be qualified with schema names.

- *`query_block_name`*: The query block to which the hint applies. If the hint includes no leading *`@query_block_name`*, the hint applies to the query block in which it occurs. For *`tbl_name@query_block_name`* syntax, the hint applies to the named table in the named query block. To assign a name to a query block, see Optimizer Hints for Naming Query Blocks.

Examples:

```
SELECT /*+ NO_BKA(t1, t2) */ t1.* FROM t1 INNER JOIN t2 INNER JOIN t3;
SELECT /*+ NO_BNL() BKA(t1) */ t1.* FROM t1 INNER JOIN t2 INNER JOIN t3;
SELECT /*+ NO_MERGE(dt) */ * FROM (SELECT * FROM t1) AS dt;
```

A table-level hint applies to tables that receive records from previous tables, not sender tables. Consider this statement:

```
SELECT /*+ BNL(t2) */ FROM t1, t2;
```

If the optimizer chooses to process `t1` first, it applies a Block Nested-Loop join to `t2` by buffering the rows from `t1` before starting to read from `t2`. If the optimizer instead chooses to process `t2` first, the hint has no effect because `t2` is a sender table.

For the `MERGE` and `NO_MERGE` hints, these precedence rules apply:

- A hint takes precedence over any optimizer heuristic that is not a technical constraint. (If providing a hint as a suggestion has no effect, the optimizer has a reason for ignoring it.)

- A hint takes precedence over the `derived_merge` flag of the `optimizer_switch` system variable.
- For view references, an `ALGORITHM={MERGE|TEMPTABLE}` clause in the view definition takes precedence over a hint specified in the query referencing the view.

Index-Level Optimizer Hints

Index-level hints affect which index-processing strategies the optimizer uses for particular tables or indexes. These hint types affect use of Index Condition Pushdown (ICP), Multi-Range Read (MRR), Index Merge, and range optimizations (see Section 8.2.1, “Optimizing SELECT Statements”).

Syntax of index-level hints:

```
hint_name([@query_block_name] tbl_name [index_name [, index_name] ...])
hint_name(tbl_name@query_block_name [index_name [, index_name] ...])
```

The syntax refers to these terms:

- *`hint_name`*: These hint names are permitted:
  - `INDEX_MERGE`, `NO_INDEX_MERGE`: Enable or disable the Index Merge access method for the specified tables or indexes. For information about this access method, see Section 8.2.1.3, “Index Merge Optimization”. These hints apply to all three Index Merge algorithms.

    The `INDEX_MERGE` hint forces the optimizer to use Index Merge for the specified table using the specified set of indexes. If no index is specified, the optimizer considers all possible index combinations and selects the least expensive one. The hint may be ignored if the index combination is inapplicable to the given statement.

    The `NO_INDEX_MERGE` hint disables Index Merge combinations that involve any of the specified indexes. If the hint specifies no indexes, Index Merge is not permitted for the table.
  - `MRR`, `NO_MRR`: Enable or disable MRR for the specified tables or indexes. MRR hints apply only to `InnoDB` and `MyISAM` tables.
  - `NO_ICP`: Disable ICP for the specified tables or indexes. By default, ICP is a candidate optimization strategy, so there is no hint for enabling it.
  - `NO_RANGE_OPTIMIZATION`: Disable index range access for the specified tables or indexes. This hint also disables Index Merge and Loose Index Scan for the tables or indexes. By default, range access is a candidate optimization strategy, so there is no hint for enabling it.

    This hint may be useful when the number of ranges may be high and range optimization would require many resources.
- *`tbl_name`*: The table to which the hint applies.`
- *`index_name`*: The name of an index in the named table. The hint applies to all indexes that it names. If the hint names no indexes, it applies to all indexes in the table.

  To refer to a primary key, use the name `PRIMARY`. To see the index names for a table, use `SHOW INDEX`.
- *`query_block_name`*: The query block to which the hint applies. If the hint includes no leading *`@query_block_name`*, the hint applies to the query block in which it occurs. For *`tbl_name@query_block_name`* syntax, the hint applies to the named table in the named query block. To assign a name to a query block, see Optimizer Hints for Naming Query Blocks.

Examples:

```
SELECT /*+ INDEX_MERGE(t1 f3, PRIMARY) */ f2 FROM t1
  WHERE f1 = 'o' AND f2 = f3 AND f3 <= 4;
SELECT /*+ MRR(t1) */ * FROM t1 WHERE f2 <= 3 AND 3 <= f3;
SELECT /*+ NO_RANGE_OPTIMIZATION(t3 PRIMARY, f2_idx) */ f1
  FROM t3 WHERE f1 > 30 AND f1 < 33;
INSERT INTO t3(f1, f2, f3)
  (SELECT /*+ NO_ICP(t2) */ t2.f1, t2.f2, t2.f3 FROM t1,t2
   WHERE t1.f1=t2.f1 AND t2.f2 BETWEEN t1.f1
   AND t1.f2 AND t2.f2 + 1 >= t1.f1 + 1);
```

Index Merge examples. Assume that table `t1` has columns `a`, `b`, `c`, and `d`; and that indexes named `i_a`, `i_b`, and `i_c` exist on `a`, `b`, and `c`, respectively:

```
SELECT /*+ INDEX_MERGE(t1 i_a, i_b, i_c)*/ * FROM t1
  WHERE a = 1 AND b = 2 AND c = 3 AND d = 4;
```

Index Merge is used for `(i_a, i_b, i_c)` in this case.

```
SELECT /*+ INDEX_MERGE(t1 i_a, i_b, i_c)*/ * FROM t1
  WHERE b = 1 AND c = 2 AND d = 3;
```

Index Merge is used for `(i_b, i_c)` in this case.

```
/*+ INDEX_MERGE(t1 i_a, i_b) NO_INDEX_MERGE(t1 i_b) */
```

`NO_INDEX_MERGE` is ignored because there is a preceding hint for the same table.

```
/*+ NO_INDEX_MERGE(t1 i_a, i_b) INDEX_MERGE(t1 i_b) */
```

`INDEX_MERGE` is ignored because there is a preceding hint for the same table.

For the `INDEX_MERGE` and `NO_INDEX_MERGE` optimizer hints, these precedence rules apply:

- If an optimizer hint is specified and is applicable, it takes precedence over the Index Merge-related flags of the `optimizer_switch` system variable.

  ```
  SET optimizer_switch='index_merge_intersection=off';
  SELECT /*+ INDEX_MERGE(t1 i_b, i_c) */ * FROM t1
  WHERE b = 1 AND c = 2 AND d = 3;
  ```

  The hint takes precedence over `optimizer_switch`. Index Merge is used for `(i_b, i_c)` in this case.

  ```
  SET optimizer_switch='index_merge_intersection=on';
  SELECT /*+ INDEX_MERGE(t1 i_b) */ * FROM t1
  WHERE b = 1 AND c = 2 AND d = 3;
  ```

  The hint specifies only one index, so it is inapplicable, and the `optimizer_switch` flag (`on`) applies. Index Merge is used if the optimizer assesses it to be cost efficient.

  ```
  SET optimizer_switch='index_merge_intersection=off';
  SELECT /*+ INDEX_MERGE(t1 i_b) */ * FROM t1
  WHERE b = 1 AND c = 2 AND d = 3;
  ```

  The hint specifies only one index, so it is inapplicable, and the `optimizer_switch` flag (`off`) applies. Index Merge is not used.

- The `USE INDEX`, `FORCE INDEX`, and `IGNORE INDEX` index hints have higher priority than the `INDEX_MERGE` and `NO_INDEX_MERGE` optimizer hints.

  ```
  ```

```
/*+ INDEX_MERGE(t1 i_a, i_b, i_c) */ ... IGNORE INDEX i_a
```

`IGNORE INDEX` takes precedence over `INDEX_MERGE`, so index `i_a` is excluded from the possible ranges for Index Merge.

```
/*+ NO_INDEX_MERGE(t1 i_a, i_b) */ ... FORCE INDEX i_a, i_b
```

Index Merge is disallowed for `i_a, i_b` because of `FORCE INDEX`, but the optimizer is forced to use either `i_a` or `i_b` for `range` or `ref` access. There are no conflicts; both hints are applicable.

- If an `IGNORE INDEX` hint names multiple indexes, those indexes are unavailable for Index Merge.
- The `FORCE INDEX` and `USE INDEX` hints make only the named indexes to be available for Index Merge.

```
SELECT /*+ INDEX_MERGE(t1 i_a, i_b, i_c) */ a FROM t1
FORCE INDEX (i_a, i_b) WHERE c = 'h' AND a = 2 AND b = 'b';
```

The Index Merge intersection access algorithm is used for `(i_a, i_b)`. The same is true if `FORCE INDEX` is changed to `USE INDEX`.

Subquery Optimizer Hints

Subquery hints affect whether to use semi-join transformations and which semi-join strategies to permit, and, when semi-joins are not used, whether to use subquery materialization or `IN`-to-`EXISTS` transformations. For more information about these optimizations, see Section 8.2.2, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions”.

Syntax of hints that affect semi-join strategies:

```
hint_name([@query_block_name] [strategy [, strategy] ...])
```

The syntax refers to these terms:

- *`hint_name`*: These hint names are permitted:
  - `SEMIJOIN`, `NO_SEMIJOIN`: Enable or disable the named semi-join strategies.
- *`strategy`*: A semi-join strategy to be enabled or disabled. These strategy names are permitted: `DUPSWEEDOUT`, `FIRSTMATCH`, `LOOSESCAN`, `MATERIALIZATION`.

  For `SEMIJOIN` hints, if no strategies are named, semi-join is used if possible based on the strategies enabled according to the `optimizer_switch` system variable. If strategies are named but inapplicable for the statement, `DUPSWEEDOUT` is used.

  For `NO_SEMIJOIN` hints, if no strategies are named, semi-join is not used. If strategies are named that rule out all applicable strategies for the statement, `DUPSWEEDOUT` is used.

If one subquery is nested within another and both are merged into a semi-join of an outer query, any specification of semi-join strategies for the innermost query are ignored. `SEMIJOIN` and `NO_SEMIJOIN` hints can still be used to enable or disable semi-join transformations for such nested subqueries.

If `DUPSWEEDOUT` is disabled, on occasion the optimizer may generate a query plan that is far from optimal. This occurs due to heuristic pruning during greedy search, which can be avoided by setting `optimizer_prune_level=0`.

Examples:

```
SELECT /*+ NO_SEMIJOIN(@subq1 FIRSTMATCH, LOOSESCAN) */ * FROM t2
  WHERE t2.a IN (SELECT /*+ QB_NAME(subq1) */ a FROM t3);
```

```
SELECT /*+ SEMIJOIN(@subq1 MATERIALIZATION, DUPSWEEDOUT) */ * FROM t2
  WHERE t2.a IN (SELECT /*+ QB_NAME(subq1) */ a FROM t3);
```

Syntax of hints that affect whether to use subquery materialization or `IN`-to-`EXISTS` transformations:

```
SUBQUERY([@query_block_name] strategy)
```

The hint name is always `SUBQUERY`.

For `SUBQUERY` hints, these *`strategy`* values are permitted: `INTOEXISTS`, `MATERIALIZATION`.

Examples:

```
SELECT id, a IN (SELECT /*+ SUBQUERY(MATERIALIZATION) */ a FROM t1) FROM t2;
SELECT * FROM t2 WHERE t2.a IN (SELECT /*+ SUBQUERY(INTOEXISTS) */ a FROM t1);
```

For semi-join and `SUBQUERY` hints, a leading *`@query_block_name`* specifies the query block to which the hint applies. If the hint includes no leading *`@query_block_name`*, the hint applies to the query block in which it occurs. To assign a name to a query block, see Optimizer Hints for Naming Query Blocks.

If a hint comment contains multiple subquery hints, the first is used. If there are other following hints of that type, they produce a warning. Following hints of other types are silently ignored.

Statement Execution Time Optimizer Hints

The `MAX_EXECUTION_TIME` hint is permitted only for `SELECT` statements. It places a limit *`N`* (a timeout value in milliseconds) on how long a statement is permitted to execute before the server terminates it:

```
MAX_EXECUTION_TIME(N)
```

Example with a timeout of 1 second (1000 milliseconds):

```
SELECT /*+ MAX_EXECUTION_TIME(1000) */ * FROM t1 INNER JOIN t2 WHERE ...
```

The `MAX_EXECUTION_TIME(N)` hint sets a statement execution timeout of *`N`* milliseconds. If this option is absent or *`N`* is 0, the statement timeout established by the `max_execution_time` system variable applies.

The `MAX_EXECUTION_TIME` hint is applicable as follows:

- For statements with multiple `SELECT` keywords, such as unions or statements with subqueries, `MAX_EXECUTION_TIME` applies to the entire statement and must appear after the first `SELECT`.
- It applies to read-only `SELECT` statements. Statements that are not read only are those that invoke a stored function that modifies data as a side effect.
- It does not apply to `SELECT` statements in stored programs and is ignored.

Optimizer Hints for Naming Query Blocks

Table-level, index-level, and subquery optimizer hints permit specific query blocks to be named as part of their argument syntax. To create these names, use the `QB_NAME` hint, which assigns a name to the query block in which it occurs:

```
QB_NAME(name)
```

`QB_NAME` hints can be used to make explicit in a clear way which query blocks other hints apply to. They also permit all non-query block name hints to be specified within a single hint comment for easier understanding of complex statements. Consider the following statement:

```
SELECT ...
  FROM (SELECT ...
  FROM (SELECT ... FROM ...)) ...
```

`QB_NAME` hints assign names to query blocks in the statement:

```
SELECT /*+ QB_NAME(qb1) */ ...
  FROM (SELECT /*+ QB_NAME(qb2) */ ...
  FROM (SELECT /*+ QB_NAME(qb3) */ ... FROM ...)) ...
```

Then other hints can use those names to refer to the appropriate query blocks:

```
SELECT /*+ QB_NAME(qb1) MRR(@qb1 t1) BKA(@qb2) NO_MRR(@qb3t1 idx1, id2) */ ...
  FROM (SELECT /*+ QB_NAME(qb2) */ ...
  FROM (SELECT /*+ QB_NAME(qb3) */ ... FROM ...)) ...
```

The resulting effect is as follows:

- `MRR(@qb1 t1)` applies to table `t1` in query block `qb1`.
- `BKA(@qb2)` applies to query block `qb2`.
- `NO_MRR(@qb3 t1 idx1, id2)` applies to indexes `idx1` and `idx2` in table `t1` in query block `qb3`.

Query block names are identifiers and follow the usual rules about what names are valid and how to quote them (see Section 9.2, “Schema Object Names”). For example, a query block name that contains spaces must be quoted, which can be done using backticks:

```
SELECT /*+ BKA(@`my hint name`) */ ...
  FROM (SELECT /*+ QB_NAME(`my hint name`) */ ...) ...
```

If the `ANSI_QUOTES` SQL mode is enabled, it is also possible to quote query block names within double quotation marks:

```
SELECT /*+ BKA(@"my hint name") */ ...
  FROM (SELECT /*+ QB_NAME("my hint name") */ ...) ...
```

8.9.3 Switchable Optimizations

The `optimizer_switch` system variable enables control over optimizer behavior. Its value is a set of flags, each of which has a value of `on` or `off` to indicate whether the corresponding optimizer behavior is enabled or disabled. This variable has global and session values and can be changed at runtime. The global default can be set at server startup.

To see the current set of optimizer flags, select the variable value:

```
mysql> SELECT @@optimizer_switch\G
*************************** 1. row ***************************
@@optimizer_switch: index_merge=on,index_merge_union=on,
                    index_merge_sort_union=on,
                    index_merge_intersection=on,
                    engine_condition_pushdown=on,
                    index_condition_pushdown=on,
                    mrr=on,mrr_cost_based=on,
                    block_nested_loop=on,batched_key_access=off,
                    materialization=on,semijoin=on,loosescan=on,
                    firstmatch=on,duplicateweedout=on,
                    subquery_materialization_cost_based=on,
                    use_index_extensions=on,
                    condition_fanout_filter=on,derived_merge=on
```

To change the value of `optimizer_switch`, assign a value consisting of a comma-separated list of one or more commands:

```
SET [GLOBAL|SESSION] optimizer_switch='command[,command]...';
```

Each *`command`* value should have one of the forms shown in the following table.

Command Syntax	Meaning
`default`	Reset every optimization to its default value
`opt_name``=default`	Set the named optimization to its default value
`opt_name``=off`	Disable the named optimization
`opt_name``=on`	Enable the named optimization

The order of the commands in the value does not matter, although the `default` command is executed first if present. Setting an *`opt_name`* flag to `default` sets it to whichever of `on` or `off` is its default value. Specifying any given *`opt_name`* more than once in the value is not permitted and causes an error. Any errors in the value cause the assignment to fail with an error, leaving the value of `optimizer_switch` unchanged.

The following table lists the permissible *`opt_name`* flag names, grouped by optimization strategy.

Optimization	Flag Name	Meaning	Default
Batched Key Access	`batched_key_access`	Controls use of BKA join algorithm	`OFF`
Block Nested-Loop	`block_nested_loop`	Controls use of BNL join algorithm	`ON`
Condition Filtering	`condition_fanout_filter`	Controls use of condition filtering	`ON`
Engine Condition Pushdown	`engine_condition_pushdown`	Controls engine condition pushdown	`ON`
Index Condition Pushdown	`index_condition_pushdown`	Controls index condition pushdown	`ON`
Index Extensions	`use_index_extensions`	Controls use of index extensions	`ON`
Index Merge	`index_merge`	Controls all Index Merge optimizations	`ON`
	`index_merge_intersection`	Controls the Index Merge Intersection Access optimization	`ON`
	`index_merge_sort_union`	Controls the Index Merge Sort-Union Access optimization	`ON`
	`index_merge_union`	Controls the Index Merge Union Access optimization	`ON`
Multi-Range Read	`mrr`	Controls the Multi-Range Read strategy	`ON`
	`mrr_cost_based`	Controls use of cost-based MRR if `mrr=on`	`ON`
Semi-join	`semijoin`	Controls all semi-join strategies	`ON`
	`firstmatch`	Controls the semi-join FirstMatch strategy	`ON`
	`loosescan`	Controls the semi-join LooseScan strategy (not to be confused with LooseScan for `GROUP BY`)	`ON`
	`duplicateweedout`	Controls the semi-join Duplicate Weedout strategy	`ON`

Optimization	Flag Name	Meaning	Default
Subquery materialization	`materialization`	Controls materialization (including semi-join materialization)	`ON`
	`subquery_materialization_cost_based`	Use cost-based materialization choice	`ON`
Derived table merging	`derived_merge`	Controls merging of derived tables and views into outer query block	`ON`

For `batched_key_access` to have any effect when set to `on`, the `mrr` flag must also be `on`. Currently, the cost estimation for MRR is too pessimistic. Hence, it is also necessary for `mrr_cost_based` to be `off` for BKA to be used.

The `semijoin`, `firstmatch`, `loosescan`, `duplicateweedout`, and `materialization` flags enable control over semi-join and subquery materialization strategies. The `semijoin` flag controls whether semi-joins are used. If it is set to `on`, the `firstmatch` and `loosescan` flags enable finer control over the permitted semi-join strategies. The `materialization` flag controls whether subquery materialization is used. If `semijoin` and `materialization` are both `on`, semi-joins also use materialization where applicable. These flags are `on` by default.

If the `duplicateweedout` semi-join strategy is disabled, it is not used unless all other applicable strategies are also disabled.

The `subquery_materialization_cost_based` flag enables control over the choice between subquery materialization and `IN`-to-`EXISTS` subquery transformation. If the flag is `on` (the default), the optimizer performs a cost-based choice between subquery materialization and `IN`-to-`EXISTS` subquery transformation if either method could be used. If the flag is `off`, the optimizer chooses subquery materialization over `IN -> EXISTS` subquery transformation.

The `derived_merge` flag controls whether the optimizer attempts to merge derived tables, view references, and common table expressions into the outer query block, assuming that no other rule prevents merging; for example, an `ALGORITHM` directive for a view takes precedence over the `derived_merge` setting. By default, the flag is `on` to enable merging. For more information, see Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

For more information about individual optimization strategies, see the following sections:

- Section 8.2.1.3, “Index Merge Optimization”
- Section 8.2.1.4, “Engine Condition Pushdown Optimization”
- Section 8.2.1.5, “Index Condition Pushdown Optimization”
- Section 8.2.1.10, “Multi-Range Read Optimization”
- Section 8.2.1.11, “Block Nested-Loop and Batched Key Access Joins”
- Section 8.2.2, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions”
- Section 8.3.9, “Use of Index Extensions”

When you assign a value to `optimizer_switch`, flags that are not mentioned keep their current values. This makes it possible to enable or disable specific optimizer behaviors in a single statement without affecting other behaviors. The statement does not depend on what other optimizer flags exist and what their values are. Suppose that all Index Merge optimizations are enabled:

```
mysql> SELECT @@optimizer_switch\G
*************************** 1. row ***************************
@@optimizer_switch: index_merge=on,index_merge_union=on,
                    index_merge_sort_union=on,
```

```
                    index_merge_intersection=on,
                    engine_condition_pushdown=on,
                    index_condition_pushdown=on,
                    mrr=on,mrr_cost_based=on,
                    block_nested_loop=on,batched_key_access=off,
                    materialization=on,semijoin=on,loosescan=on,
                    firstmatch=on,
                    subquery_materialization_cost_based=on,
                    use_index_extensions=on,
                    condition_fanout_filter=on
```

If the server is using the Index Merge Union or Index Merge Sort-Union access methods for certain queries and you want to check whether the optimizer will perform better without them, set the variable value like this:

```
mysql> SET optimizer_switch='index_merge_union=off,index_merge_sort_union=off';

mysql> SELECT @@optimizer_switch\G
*************************** 1. row ***************************
@@optimizer_switch: index_merge=on,index_merge_union=off,
                    index_merge_sort_union=off,
                    index_merge_intersection=on,
                    engine_condition_pushdown=on,
                    index_condition_pushdown=on,
                    mrr=on,mrr_cost_based=on,
                    block_nested_loop=on,batched_key_access=off,
                    materialization=on,semijoin=on,loosescan=on,
                    firstmatch=on,
                    subquery_materialization_cost_based=on,
                    use_index_extensions=on,
                    condition_fanout_filter=on
```

8.9.4 Index Hints

Index hints give the optimizer information about how to choose indexes during query processing. Index hints, described here, differ from optimizer hints, described in Section 8.9.2, “Optimizer Hints”. Index and optimizer hints may be used separately or together.

Index hints are specified following a table name. (For the general syntax for specifying tables in a `SELECT` statement, see Section 13.2.10.2, “JOIN Syntax”.) The syntax for referring to an individual table, including index hints, looks like this:

```
tbl_name [[AS] alias] [index_hint_list]

index_hint_list:
    index_hint [, index_hint] ...

index_hint:
    USE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] ([index_list])
  | IGNORE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] (index_list)
  | FORCE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] (index_list)

index_list:
    index_name [, index_name] ...
```

The `USE INDEX (index_list)` hint tells MySQL to use only one of the named indexes to find rows in the table. The alternative syntax `IGNORE INDEX (index_list)` tells MySQL to not use some particular index or indexes. These hints are useful if `EXPLAIN` shows that MySQL is using the wrong index from the list of possible indexes.

The `FORCE INDEX` hint acts like `USE INDEX (index_list)`, with the addition that a table scan is assumed to be *very* expensive. In other words, a table scan is used only if there is no way to use one of the named indexes to find rows in the table.

Each hint requires the names of *indexes*, not the names of columns. To refer to a primary key, use the name `PRIMARY`. To see the index names for a table, use `SHOW INDEX`.

An *`index_name`* value need not be a full index name. It can be an unambiguous prefix of an index name. If a prefix is ambiguous, an error occurs.

Examples:

```
SELECT * FROM table1 USE INDEX (col1_index,col2_index)
  WHERE col1=1 AND col2=2 AND col3=3;

SELECT * FROM table1 IGNORE INDEX (col3_index)
  WHERE col1=1 AND col2=2 AND col3=3;
```

The syntax for index hints has the following characteristics:

- It is syntactically valid to omit *`index_list`* for `USE INDEX`, which means “use no indexes.” Omitting *`index_list`* for `FORCE INDEX` or `IGNORE INDEX` is a syntax error.
- You can specify the scope of an index hint by adding a `FOR` clause to the hint. This provides more fine-grained control over the optimizer's selection of an execution plan for various phases of query processing. To affect only the indexes used when MySQL decides how to find rows in the table and how to process joins, use `FOR JOIN`. To influence index usage for sorting or grouping rows, use `FOR ORDER BY` or `FOR GROUP BY`.
- You can specify multiple index hints:

  ```
  SELECT * FROM t1 USE INDEX (i1) IGNORE INDEX FOR ORDER BY (i2) ORDER BY a;
  ```

  It is not an error to name the same index in several hints (even within the same hint):

  ```
  SELECT * FROM t1 USE INDEX (i1) USE INDEX (i1,i1);
  ```

  However, it is an error to mix `USE INDEX` and `FORCE INDEX` for the same table:

  ```
  SELECT * FROM t1 USE INDEX FOR JOIN (i1) FORCE INDEX FOR JOIN (i2);
  ```

If an index hint includes no `FOR` clause, the scope of the hint is to apply to all parts of the statement. For example, this hint:

```
IGNORE INDEX (i1)
```

is equivalent to this combination of hints:

```
IGNORE INDEX FOR JOIN (i1)
IGNORE INDEX FOR ORDER BY (i1)
IGNORE INDEX FOR GROUP BY (i1)
```

In MySQL 5.0, hint scope with no `FOR` clause was to apply only to row retrieval. To cause the server to use this older behavior when no `FOR` clause is present, enable the `old` system variable at server startup. Take care about enabling this variable in a replication setup. With statement-based binary logging, having different modes for the master and slaves might lead to replication errors.

When index hints are processed, they are collected in a single list by type (`USE`, `FORCE`, `IGNORE`) and by scope (`FOR JOIN`, `FOR ORDER BY`, `FOR GROUP BY`). For example:

```
SELECT * FROM t1
  USE INDEX () IGNORE INDEX (i2) USE INDEX (i1) USE INDEX (i2);
```

is equivalent to:

```
SELECT * FROM t1
   USE INDEX (i1,i2) IGNORE INDEX (i2);
```

The index hints then are applied for each scope in the following order:

1. `{USE|FORCE} INDEX` is applied if present. (If not, the optimizer-determined set of indexes is used.)
2. `IGNORE INDEX` is applied over the result of the previous step. For example, the following two queries are equivalent:

```
SELECT * FROM t1 USE INDEX (i1) IGNORE INDEX (i2) USE INDEX (i2);

SELECT * FROM t1 USE INDEX (i1);
```

For `FULLTEXT` searches, index hints work as follows:

- For natural language mode searches, index hints are silently ignored. For example, `IGNORE INDEX(i1)` is ignored with no warning and the index is still used.
- For boolean mode searches, index hints with `FOR ORDER BY` or `FOR GROUP BY` are silently ignored. Index hints with `FOR JOIN` or no `FOR` modifier are honored. In contrast to how hints apply for non-`FULLTEXT` searches, the hint is used for all phases of query execution (finding rows and retrieval, grouping, and ordering). This is true even if the hint is given for a non-`FULLTEXT` index.

  For example, the following two queries are equivalent:

```
SELECT * FROM t
  USE INDEX (index1)
  IGNORE INDEX (index1) FOR ORDER BY
  IGNORE INDEX (index1) FOR GROUP BY
  WHERE ... IN BOOLEAN MODE ... ;

SELECT * FROM t
  USE INDEX (index1)
  WHERE ... IN BOOLEAN MODE ... ;
```

8.9.5 The Optimizer Cost Model

To generate execution plans, the optimizer uses a cost model that is based on estimates of the cost of various operations that occur during query execution. The optimizer has a set of compiled-in default “cost constants” available to it to make decisions regarding execution plans.

The optimizer also has a database of cost estimates to use during execution plan construction. These estimates are stored in the `server_cost` and `engine_cost` tables in the `mysql` system database and are configurable at any time. The intent of these tables is to make it possible to easily adjust the cost estimates that the optimizer uses when it attempts to arrive at query execution plans.

- Cost Model General Operation
- The Cost Model Database
- Making Changes to the Cost Model Database

Cost Model General Operation

The configurable optimizer cost model works like this:

- The server reads the cost model tables into memory at startup and uses the in-memory values at runtime. Any non-`NULL` cost estimate specified in the tables takes precedence over the corresponding compiled-in default cost constant. Any `NULL` estimate indicates to the optimizer to use the compiled-in default.

- At runtime, the server may reread the cost tables. This occurs when a storage engine is dynamically loaded or when a `FLUSH OPTIMIZER_COSTS` statement is executed.
- Cost tables enable server administrators to easily adjust cost estimates by changing entries in the tables. It is also easy to revert to a default by setting an entry's cost to `NULL`. The optimizer uses the in-memory cost values, so changes to the tables should be followed by `FLUSH OPTIMIZER_COSTS` to take effect.
- The in-memory cost estimates that are current when a client session begins apply throughout that session until it ends. In particular, if the server rereads the cost tables, any changed estimates apply only to subsequently started sessions. Existing sessions are unaffected.
- Cost tables are specific to a given server instance. The server does not replicate cost table changes to replication slaves.

The Cost Model Database

The optimizer cost model database consists of two tables in the `mysql` system database that contain cost estimate information for operations that occur during query execution:

- `server_cost`: Optimizer cost estimates for general server operations
- `engine_cost`: Optimizer cost estimates for operations specific to particular storage engines

The `server_cost` table contains these columns:

- `cost_name`

  The name of a cost estimate used in the cost model. The name is not case sensitive. If the server does not recognize the cost name when it reads this table, it writes a warning to the error log.

- `cost_value`

  The cost estimate value. If the value is non-`NULL`, the server uses it as the cost. Otherwise, it uses the default estimate (the compiled-in value). DBAs can change a cost estimate by updating this column. If the server finds that the cost value is invalid (nonpositive) when it reads this table, it writes a warning to the error log.

  To override a default cost estimate (for an entry that specifies `NULL`), set the cost to a non-`NULL` value. To revert to the default, set the value to `NULL`. Then execute `FLUSH OPTIMIZER_COSTS` to tell the server to reread the cost tables.

- `last_update`

  The time of the last row update.

- `comment`

  A descriptive comment associated with the cost estimate. DBAs can use this column to provide information about why a cost estimate row stores a particular value.

- `default_value`

  The default (compiled-in) value for the cost estimate. This column is a read-only generated column that retains its value even if the associated cost estimate is changed. For rows added to the table at runtime, the value of this column is `NULL`.

The primary key for the `server_cost` table is the `cost_name` column, so it is not possible to create multiple entries for any cost estimate.

The server recognizes these `cost_name` values for the `server_cost` table:

- `disk_temptable_create_cost`, `disk_temptable_row_cost`

The cost estimates for internally created temporary tables stored in a disk-based storage engine (either `InnoDB` or `MyISAM`). Increasing these values increases the cost estimate of using internal temporary tables and makes the optimizer prefer query plans with less use of them. For information about such tables, see Section 8.4.4, “Internal Temporary Table Use in MySQL”.

The larger default values for these disk parameters compared to the default values for the corresponding memory parameters (`memory_temptable_create_cost`, `memory_temptable_row_cost`) reflects the greater cost of processing disk-based tables.

- `key_compare_cost`

  The cost of comparing record keys. Increasing this value causes a query plan that compares many keys to become more expensive. For example, a query plan that performs a `filesort` becomes relatively more expensive compared to a query plan that avoids sorting by using an index.

- `memory_temptable_create_cost`, `memory_temptable_row_cost`

  The cost estimates for internally created temporary tables stored in the `MEMORY` storage engine. Increasing these values increases the cost estimate of using internal temporary tables and makes the optimizer prefer query plans with less use of them. For information about such tables, see Section 8.4.4, “Internal Temporary Table Use in MySQL”.

  The smaller default values for these memory parameters compared to the default values for the corresponding disk parameters (`disk_temptable_create_cost`, `disk_temptable_row_cost`) reflects the lesser cost of processing memory-based tables.

- `row_evaluate_cost`

  The cost of evaluating record conditions. Increasing this value causes a query plan that examines many rows to become more expensive compared to a query plan that examines fewer rows. For example, a table scan becomes relatively more expensive compared to a range scan that reads fewer rows.

The `engine_cost` table contains these columns:

- `engine_name`

  The name of the storage engine to which this cost estimate applies. The name is not case sensitive. If the value is `default`, it applies to all storage engines that have no named entry of their own. If the server does not recognize the engine name when it reads this table, it writes a warning to the error log.

- `device_type`

  The device type to which this cost estimate applies. The column is intended for specifying different cost estimates for different storage device types, such as hard disk drives versus solid state drives. Currently, this information is not used and 0 is the only permitted value.

- `cost_name`

  Same as in the `server_cost` table.

- `cost_value`

  Same as in the `server_cost` table.

- `last_update`

  Same as in the `server_cost` table.

- `comment`

Same as in the `server_cost` table.

- `default_value`

  The default (compiled-in) value for the cost estimate. This column is a read-only generated column that retains its value even if the associated cost estimate is changed. For rows added to the table at runtime, the value of this column is `NULL`, with the exception that if the row has the same `cost_name` value as one of the original rows, the `default_value` column will have the same value as that row.

The primary key for the `engine_cost` table is a tuple comprising the (`cost_name`, `engine_name`, `device_type`) columns, so it is not possible to create multiple entries for any combination of values in those columns.

The server recognizes these `cost_name` values for the `engine_cost` table:

- `io_block_read_cost`

  The cost of reading an index or data block from disk. Increasing this value causes a query plan that reads many disk blocks to become more expensive compared to a query plan that reads fewer disk blocks. For example, a table scan becomes relatively more expensive compared to a range scan that reads fewer blocks.

- `memory_block_read_cost`

  Similar to `io_block_read_cost`, but represents the cost of reading an index or data block from an in-memory database buffer.

If the `io_block_read_cost` and `memory_block_read_cost` values differ, the execution plan may change between two runs of the same query. Suppose that the cost for memory access is less than the cost for disk access. In that case, at server startup before data has been read into the buffer pool, you may get a different plan than after the query has been run because then the data will be in memory.

Making Changes to the Cost Model Database

For DBAs who wish to change the cost model parameters from their defaults, try doubling or halving the value and measuring the effect.

Changes to the `io_block_read_cost` and `memory_block_read_cost` parameters are most likely to yield worthwhile results. These parameter values enable cost models for data access methods to take into account the costs of reading information from different sources; that is, the cost of reading information from disk versus reading information already in a memory buffer. For example, all other things being equal, setting `io_block_read_cost` to a value larger than `memory_block_read_cost` causes the optimizer to prefer query plans that read information already held in memory to plans that must read from disk.

This example shows how to change the default value for `io_block_read_cost`:

```
UPDATE mysql.engine_cost
  SET cost_value = 2.0
  WHERE cost_name = 'io_block_read_cost';
FLUSH OPTIMIZER_COSTS;
```

This example shows how to change the value of `io_block_read_cost` only for the `InnoDB` storage engine:

```
INSERT INTO mysql.engine_cost
  VALUES ('InnoDB', 0, 'io_block_read_cost', 3.0,
  CURRENT_TIMESTAMP, 'Using a slower disk for InnoDB');
FLUSH OPTIMIZER_COSTS;
```

8.9.6 Optimizer Statistics

As of MySQL 8.0.2, the `column_statistics` data dictionary table stores histogram statistics about column values. To perform histogram management, use the `ANALYZE TABLE` statement; see Section 13.7.2.1, “ANALYZE TABLE Syntax”.

Note

Currently, the optimizer does not yet consult the `column_statistics` table in the course of query execution plan construction.

The `column_statistics` table has these characteristics:

- The table contains statistics for columns of all data types except geometry types (spatial data) and `JSON`.
- The table makes column statistics persistent so that they need not be created each time the server starts.
- The server performs updates to the table; users do not.

The `column_statistics` table is not directly accessible by users because it is part of the data dictionary. Histogram information is available using `INFORMATION_SCHEMA.COLUMN_STATISTICS`, which is implemented as a view on the data dictionary table. `COLUMN_STATISTICS` has these columns:

- `SCHEMA_NAME`, `TABLE_NAME`, `COLUMN_NAME`: The names of the schema, table, and column for which the statistics apply.
- `HISTOGRAM`: A `JSON` value describing the column statistics, stored as a histogram.

Column histograms contain buckets for parts of the range of values stored in the column. Histograms are `JSON` objects to permit flexibility in the way column statistics are represented. Here is a sample histogram object:

```
{
  "buckets": [
    [
      1,
      0.3333333333333333
    ],
    [
      2,
      0.6666666666666666
    ],
    [
      3,
      1
    ]
  ],
  "null-values": 0,
  "last-updated": "2017-03-24 13:32:40.000000",
  "sampling-rate": 1,
  "histogram-type": "singleton",
  "number-of-buckets-specified": 128,
  "data-type": "int",
  "charset-id": 8
}
```

Histogram objects have these keys:

- `buckets`: The histogram buckets. Bucket structure depends on the histogram type.

  For `singleton` histograms, buckets contain two values:

  - Value 1: The value for the bucket. The type depends on the column data type.

    - Value 2: A double representing the cumulative frequency for the value. For example, .25 and .75 indicate that 25% and 75% of the values in the column are less than or equal to the bucket value.

    For `equi-height` histograms, buckets contain four values:

    - Values 1, 2: The lower and upper inclusive values for the bucket. The type depends on the column data type.
    - Value 3: A double representing the cumulative frequency for the value. For example, .25 and .75 indicate that 25% and 75% of the values in the column are less than or equal to the bucket upper value.
    - Value 4: The number of distinct values in the range from the bucket lower value to its upper value.

- `null-values`: A number between 0.0 and 1.0 indicating the fraction of column values that are SQL `NULL` values. If 0, the column contains no `NULL` values.
- `last-updated`: When the histogram was generated, as a UTC value in `YYYY-MM-DD HH:MM:SS.000000` format (the fractional seconds part is always zero).
- `sampling-rate`: A number between 0.0 and 1.0 indicating the fraction of data that was sampled to create the histogram. A value of 1 means that all of the data was read (no sampling).
- `histogram-type`: The histogram type:
    - `singleton`: One bucket represents one single value in the column. This histogram type is created when the number of distinct values in the column is less than or equal to the number of buckets specified in the `ANALYZE TABLE` statement that generated the histogram.
    - `equi-height`: One bucket represents a range of values. This histogram type is created when the number of distinct values in the column is greater than the number of buckets specified in the `ANALYZE TABLE` statement that generated the histogram.
- `number-of-buckets-specified`: The number of buckets specified in the `ANALYZE TABLE` statement that generated the histogram.
- `data-type`: The type of data this histogram contains. This is needed when reading and parsing histograms from persistent storage into memory. The value is one of `int`, `uint` (unsigned integer), `double`, `decimal`, `datetime`, or `string` (includes character and binary strings).
- `charset-id`: The collation ID for the histogram data. It is mostly meaningful when the `data-type` value is `string`. Values correspond to `ID` column values in the `INFORMATION_SCHEMA.COLLATIONS` table.

8.10 Buffering and Caching

MySQL uses several strategies that cache information in memory buffers to increase performance.

8.10.1 InnoDB Buffer Pool Optimization

`InnoDB` maintains a storage area called the buffer pool for caching data and indexes in memory. Knowing how the `InnoDB` buffer pool works, and taking advantage of it to keep frequently accessed data in memory, is an important aspect of MySQL tuning.

For an explanation of the inner workings of the `InnoDB` buffer pool, an overview of its LRU replacement algorithm, and general configuration information, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

For additional `InnoDB` buffer pool configuration and tuning information, see these sections:

- Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”

- Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing”
- Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”
- Section 15.6.3.3, “Configuring Multiple Buffer Pool Instances”
- Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”
- Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”
- Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”

8.10.2 The MyISAM Key Cache

To minimize disk I/O, the `MyISAM` storage engine exploits a strategy that is used by many database management systems. It employs a cache mechanism to keep the most frequently accessed table blocks in memory:

- For index blocks, a special structure called the *key cache* (or *key buffer*) is maintained. The structure contains a number of block buffers where the most-used index blocks are placed.
- For data blocks, MySQL uses no special cache. Instead it relies on the native operating system file system cache.

This section first describes the basic operation of the `MyISAM` key cache. Then it discusses features that improve key cache performance and that enable you to better control cache operation:

- Multiple sessions can access the cache concurrently.
- You can set up multiple key caches and assign table indexes to specific caches.

To control the size of the key cache, use the `key_buffer_size` system variable. If this variable is set equal to zero, no key cache is used. The key cache also is not used if the `key_buffer_size` value is too small to allocate the minimal number of block buffers (8).

When the key cache is not operational, index files are accessed using only the native file system buffering provided by the operating system. (In other words, table index blocks are accessed using the same strategy as that employed for table data blocks.)

An index block is a contiguous unit of access to the `MyISAM` index files. Usually the size of an index block is equal to the size of nodes of the index B-tree. (Indexes are represented on disk using a B-tree data structure. Nodes at the bottom of the tree are leaf nodes. Nodes above the leaf nodes are nonleaf nodes.)

All block buffers in a key cache structure are the same size. This size can be equal to, greater than, or less than the size of a table index block. Usually one these two values is a multiple of the other.

When data from any table index block must be accessed, the server first checks whether it is available in some block buffer of the key cache. If it is, the server accesses data in the key cache rather than on disk. That is, it reads from the cache or writes into it rather than reading from or writing to disk. Otherwise, the server chooses a cache block buffer containing a different table index block (or blocks) and replaces the data there by a copy of required table index block. As soon as the new index block is in the cache, the index data can be accessed.

If it happens that a block selected for replacement has been modified, the block is considered “dirty.” In this case, prior to being replaced, its contents are flushed to the table index from which it came.

Usually the server follows an *LRU (Least Recently Used)* strategy: When choosing a block for replacement, it selects the least recently used index block. To make this choice easier, the key cache module maintains all used blocks in a special list (*LRU chain*) ordered by time of use. When a block is accessed, it is the most recently used and is placed at the end of the list. When blocks need to be replaced, blocks at the beginning of the list are the least recently used and become the first candidates for eviction.

The InnoDB storage engine also uses an LRU algorithm, to manage its buffer pool. See Section 15.6.3.1, “The InnoDB Buffer Pool”.

8.10.2.1 Shared Key Cache Access

Threads can access key cache buffers simultaneously, subject to the following conditions:

- A buffer that is not being updated can be accessed by multiple sessions.
- A buffer that is being updated causes sessions that need to use it to wait until the update is complete.
- Multiple sessions can initiate requests that result in cache block replacements, as long as they do not interfere with each other (that is, as long as they need different index blocks, and thus cause different cache blocks to be replaced).

Shared access to the key cache enables the server to improve throughput significantly.

8.10.2.2 Multiple Key Caches

Shared access to the key cache improves performance but does not eliminate contention among sessions entirely. They still compete for control structures that manage access to the key cache buffers. To reduce key cache access contention further, MySQL also provides multiple key caches. This feature enables you to assign different table indexes to different key caches.

Where there are multiple key caches, the server must know which cache to use when processing queries for a given MyISAM table. By default, all MyISAM table indexes are cached in the default key cache. To assign table indexes to a specific key cache, use the CACHE INDEX statement (see Section 13.7.6.2, “CACHE INDEX Syntax”). For example, the following statement assigns indexes from the tables t1, t2, and t3 to the key cache named hot_cache:

```
mysql> CACHE INDEX t1, t2, t3 IN hot_cache;
+---------+--------------------+----------+----------+
| Table   | Op                 | Msg_type | Msg_text |
+---------+--------------------+----------+----------+
| test.t1 | assign_to_keycache | status   | OK       |
| test.t2 | assign_to_keycache | status   | OK       |
| test.t3 | assign_to_keycache | status   | OK       |
+---------+--------------------+----------+----------+
```

The key cache referred to in a CACHE INDEX statement can be created by setting its size with a SET GLOBAL parameter setting statement or by using server startup options. For example:

```
mysql> SET GLOBAL keycache1.key_buffer_size=128*1024;
```

To destroy a key cache, set its size to zero:

```
mysql> SET GLOBAL keycache1.key_buffer_size=0;
```

You cannot destroy the default key cache. Any attempt to do this is ignored:

```
mysql> SET GLOBAL key_buffer_size = 0;

mysql> SHOW VARIABLES LIKE 'key_buffer_size';
+-----------------+---------+
| Variable_name   | Value   |
+-----------------+---------+
| key_buffer_size | 8384512 |
+-----------------+---------+
```

Key cache variables are structured system variables that have a name and components. For keycache1.key_buffer_size, keycache1 is the cache variable name and key_buffer_size

is the cache component. See Section 5.1.6.1, “Structured System Variables”, for a description of the syntax used for referring to structured key cache system variables.

By default, table indexes are assigned to the main (default) key cache created at the server startup. When a key cache is destroyed, all indexes assigned to it are reassigned to the default key cache.

For a busy server, you can use a strategy that involves three key caches:

- A “hot” key cache that takes up 20% of the space allocated for all key caches. Use this for tables that are heavily used for searches but that are not updated.
- A “cold” key cache that takes up 20% of the space allocated for all key caches. Use this cache for medium-sized, intensively modified tables, such as temporary tables.
- A “warm” key cache that takes up 60% of the key cache space. Employ this as the default key cache, to be used by default for all other tables.

One reason the use of three key caches is beneficial is that access to one key cache structure does not block access to the others. Statements that access tables assigned to one cache do not compete with statements that access tables assigned to another cache. Performance gains occur for other reasons as well:

- The hot cache is used only for retrieval queries, so its contents are never modified. Consequently, whenever an index block needs to be pulled in from disk, the contents of the cache block chosen for replacement need not be flushed first.
- For an index assigned to the hot cache, if there are no queries requiring an index scan, there is a high probability that the index blocks corresponding to nonleaf nodes of the index B-tree remain in the cache.
- An update operation most frequently executed for temporary tables is performed much faster when the updated node is in the cache and need not be read in from disk first. If the size of the indexes of the temporary tables are comparable with the size of cold key cache, the probability is very high that the updated node is in the cache.

The `CACHE INDEX` statement sets up an association between a table and a key cache, but the association is lost each time the server restarts. If you want the association to take effect each time the server starts, one way to accomplish this is to use an option file: Include variable settings that configure your key caches, and an `init-file` option that names a file containing `CACHE INDEX` statements to be executed. For example:

```
key_buffer_size = 4G
hot_cache.key_buffer_size = 2G
cold_cache.key_buffer_size = 2G
init_file=/path/to/data-directory/mysqld_init.sql
```

The statements in `mysqld_init.sql` are executed each time the server starts. The file should contain one SQL statement per line. The following example assigns several tables each to `hot_cache` and `cold_cache`:

```
CACHE INDEX db1.t1, db1.t2, db2.t3 IN hot_cache
CACHE INDEX db1.t4, db2.t5, db2.t6 IN cold_cache
```

8.10.2.3 Midpoint Insertion Strategy

By default, the key cache management system uses a simple LRU strategy for choosing key cache blocks to be evicted, but it also supports a more sophisticated method called the *midpoint insertion strategy.*

When using the midpoint insertion strategy, the LRU chain is divided into two parts: a hot sublist and a warm sublist. The division point between two parts is not fixed, but the key cache

management system takes care that the warm part is not “too short,” always containing at least `key_cache_division_limit` percent of the key cache blocks. `key_cache_division_limit` is a component of structured key cache variables, so its value is a parameter that can be set per cache.

When an index block is read from a table into the key cache, it is placed at the end of the warm sublist. After a certain number of hits (accesses of the block), it is promoted to the hot sublist. At present, the number of hits required to promote a block (3) is the same for all index blocks.

A block promoted into the hot sublist is placed at the end of the list. The block then circulates within this sublist. If the block stays at the beginning of the sublist for a long enough time, it is demoted to the warm sublist. This time is determined by the value of the `key_cache_age_threshold` component of the key cache.

The threshold value prescribes that, for a key cache containing *`N`* blocks, the block at the beginning of the hot sublist not accessed within the last *`N`* `* key_cache_age_threshold / 100` hits is to be moved to the beginning of the warm sublist. It then becomes the first candidate for eviction, because blocks for replacement always are taken from the beginning of the warm sublist.

The midpoint insertion strategy enables you to keep more-valued blocks always in the cache. If you prefer to use the plain LRU strategy, leave the `key_cache_division_limit` value set to its default of 100.

The midpoint insertion strategy helps to improve performance when execution of a query that requires an index scan effectively pushes out of the cache all the index blocks corresponding to valuable high-level B-tree nodes. To avoid this, you must use a midpoint insertion strategy with the `key_cache_division_limit` set to much less than 100. Then valuable frequently hit nodes are preserved in the hot sublist during an index scan operation as well.

8.10.2.4 Index Preloading

If there are enough blocks in a key cache to hold blocks of an entire index, or at least the blocks corresponding to its nonleaf nodes, it makes sense to preload the key cache with index blocks before starting to use it. Preloading enables you to put the table index blocks into a key cache buffer in the most efficient way: by reading the index blocks from disk sequentially.

Without preloading, the blocks are still placed into the key cache as needed by queries. Although the blocks will stay in the cache, because there are enough buffers for all of them, they are fetched from disk in random order, and not sequentially.

To preload an index into a cache, use the `LOAD INDEX INTO CACHE` statement. For example, the following statement preloads nodes (index blocks) of indexes of the tables `t1` and `t2`:

```
mysql> LOAD INDEX INTO CACHE t1, t2 IGNORE LEAVES;
+---------+--------------+----------+----------+
| Table   | Op           | Msg_type | Msg_text |
+---------+--------------+----------+----------+
| test.t1 | preload_keys | status   | OK       |
| test.t2 | preload_keys | status   | OK       |
+---------+--------------+----------+----------+
```

The `IGNORE LEAVES` modifier causes only blocks for the nonleaf nodes of the index to be preloaded. Thus, the statement shown preloads all index blocks from `t1`, but only blocks for the nonleaf nodes from `t2`.

If an index has been assigned to a key cache using a `CACHE INDEX` statement, preloading places index blocks into that cache. Otherwise, the index is loaded into the default key cache.

8.10.2.5 Key Cache Block Size

It is possible to specify the size of the block buffers for an individual key cache using the `key_cache_block_size` variable. This permits tuning of the performance of I/O operations for index files.

The best performance for I/O operations is achieved when the size of read buffers is equal to the size of the native operating system I/O buffers. But setting the size of key nodes equal to the size of the I/O buffer does not always ensure the best overall performance. When reading the big leaf nodes, the server pulls in a lot of unnecessary data, effectively preventing reading other leaf nodes.

To control the size of blocks in the `.MYI` index file of `MyISAM` tables, use the `--myisam-block-size` option at server startup.

8.10.2.6 Restructuring a Key Cache

A key cache can be restructured at any time by updating its parameter values. For example:

```
mysql> SET GLOBAL cold_cache.key_buffer_size=4*1024*1024;
```

If you assign to either the `key_buffer_size` or `key_cache_block_size` key cache component a value that differs from the component's current value, the server destroys the cache's old structure and creates a new one based on the new values. If the cache contains any dirty blocks, the server saves them to disk before destroying and re-creating the cache. Restructuring does not occur if you change other key cache parameters.

When restructuring a key cache, the server first flushes the contents of any dirty buffers to disk. After that, the cache contents become unavailable. However, restructuring does not block queries that need to use indexes assigned to the cache. Instead, the server directly accesses the table indexes using native file system caching. File system caching is not as efficient as using a key cache, so although queries execute, a slowdown can be anticipated. After the cache has been restructured, it becomes available again for caching indexes assigned to it, and the use of file system caching for the indexes ceases.

8.10.3 The MySQL Query Cache

The query cache stores the text of a `SELECT` statement together with the corresponding result that was sent to the client. If an identical statement is received later, the server retrieves the results from the query cache rather than parsing and executing the statement again. The query cache is shared among sessions, so a result set generated by one client can be sent in response to the same query issued by another client.

The query cache can be useful in an environment where you have tables that do not change very often and for which the server receives many identical queries. This is a typical situation for many Web servers that generate many dynamic pages based on database content.

The query cache does not return stale data. When tables are modified, any relevant entries in the query cache are flushed.

Note

The query cache does not work in an environment where you have multiple `mysqld` servers updating the same `MyISAM` tables.

The query cache is used for prepared statements under the conditions described in Section 8.10.3.1, “How the Query Cache Operates”.

Note

The query cache is not supported for partitioned tables, and is automatically disabled for queries involving partitioned tables. The query cache cannot be enabled for such queries.

Some performance data for the query cache follows. These results were generated by running the MySQL benchmark suite on a Linux Alpha 2×500MHz system with 2GB RAM and a 64MB query cache.

- If all the queries you are performing are simple (such as selecting a row from a table with one row), but still differ so that the queries cannot be cached, the overhead for having the query cache active is 13%. This could be regarded as the worst case scenario. In real life, queries tend to be much more complicated, so the overhead normally is significantly lower.

- Searches for a single row in a single-row table are 238% faster with the query cache than without it. This can be regarded as close to the minimum speedup to be expected for a query that is cached.

To disable the query cache at server startup, set the `query_cache_size` system variable to 0. By disabling the query cache code, there is no noticeable overhead.

The query cache offers the potential for substantial performance improvement, but do not assume that it will do so under all circumstances. With some query cache configurations or server workloads, you might actually see a performance decrease:

- Be cautious about sizing the query cache excessively large, which increases the overhead required to maintain the cache, possibly beyond the benefit of enabling it. Sizes in tens of megabytes are usually beneficial. Sizes in the hundreds of megabytes might not be.

- Server workload has a significant effect on query cache efficiency. A query mix consisting almost entirely of a fixed set of `SELECT` statements is much more likely to benefit from enabling the cache than a mix in which frequent `INSERT` statements cause continual invalidation of results in the cache. In some cases, a workaround is to use the `SQL_NO_CACHE` option to prevent results from even entering the cache for `SELECT` statements that use frequently modified tables. (See Section 8.10.3.2, “Query Cache SELECT Options”.)

To verify that enabling the query cache is beneficial, test the operation of your MySQL server with the cache enabled and disabled. Then retest periodically because query cache efficiency may change as server workload changes.

8.10.3.1 How the Query Cache Operates

This section describes how the query cache works when it is operational. Section 8.10.3.3, “Query Cache Configuration”, describes how to control whether it is operational.

Incoming queries are compared to those in the query cache before parsing, so the following two queries are regarded as different by the query cache:

```
SELECT * FROM tbl_name
Select * from tbl_name
```

Queries must be *exactly* the same (byte for byte) to be seen as identical. In addition, query strings that are identical may be treated as different for other reasons. Queries that use different databases, different protocol versions, or different default character sets are considered different queries and are cached separately.

The cache is not used for queries of the following types:

- Queries that are a subquery of an outer query

- Queries executed within the body of a stored function, trigger, or event

Before a query result is fetched from the query cache, MySQL checks whether the user has `SELECT` privilege for all databases and tables involved. If this is not the case, the cached result is not used.

If a query result is returned from query cache, the server increments the `Qcache_hits` status variable, not `Com_select`. See Section 8.10.3.4, “Query Cache Status and Maintenance”.

If a table changes, all cached queries that use the table become invalid and are removed from the cache. This includes queries that use `MERGE` tables that map to the changed table. A table can be

changed by many types of statements, such as `INSERT`, `UPDATE`, `DELETE`, `TRUNCATE TABLE`, `ALTER TABLE`, `DROP TABLE`, or `DROP DATABASE`.

The query cache also works within transactions when using `InnoDB` tables.

The result from a `SELECT` query on a view is cached.

The query cache works for `SELECT SQL_CALC_FOUND_ROWS ...` queries and stores a value that is returned by a following `SELECT FOUND_ROWS()` query. `FOUND_ROWS()` returns the correct value even if the preceding query was fetched from the cache because the number of found rows is also stored in the cache. The `SELECT FOUND_ROWS()` query itself cannot be cached.

Prepared statements that are issued using the binary protocol using `mysql_stmt_prepare()` and `mysql_stmt_execute()` (see Section 27.7.8, “C API Prepared Statements”), are subject to limitations on caching. Comparison with statements in the query cache is based on the text of the statement after expansion of `?` parameter markers. The statement is compared only with other cached statements that were executed using the binary protocol. That is, for query cache purposes, prepared statements issued using the binary protocol are distinct from prepared statements issued using the text protocol (see Section 13.5, “Prepared SQL Statement Syntax”).

A query cannot be cached if it contains any of the functions shown in the following table.

`AES_DECRYPT()`	`AES_ENCRYPT()`	`BENCHMARK()`
`CONNECTION_ID()`	`CONVERT_TZ()`	`CURDATE()`
`CURRENT_DATE()`	`CURRENT_TIME()`	`CURRENT_TIMESTAMP()`
`CURRENT_USER()`	`CURTIME()`	`DATABASE()`
`ENCRYPT()` with one parameter	`FOUND_ROWS()`	`GET_LOCK()`
`IS_FREE_LOCK()`	`IS_USED_LOCK()`	`LAST_INSERT_ID()`
`LOAD_FILE()`	`MASTER_POS_WAIT()`	`NOW()`
`PASSWORD()`	`RAND()`	`RANDOM_BYTES()`
`RELEASE_ALL_LOCKS()`	`RELEASE_LOCK()`	`SLEEP()`
`SYSDATE()`	`UNIX_TIMESTAMP()` with no parameters	`USER()`
`UUID()`	`UUID_SHORT()`	

A query also is not cached under these conditions:

- It refers to user-defined functions (UDFs) or stored functions.
- It refers to user variables or local stored program variables.
- It refers to tables in the `mysql`, `INFORMATION_SCHEMA`, or `performance_schema` database.
- It refers to any partitioned tables.
- It is of any of the following forms:

```
SELECT ... FOR SHARE (or SELECT ... LOCK IN SHARE MODE)
SELECT ... FOR UPDATE
SELECT ... INTO OUTFILE ...
SELECT ... INTO DUMPFILE ...
SELECT * FROM ... WHERE autoincrement_col IS NULL
```

  The last form is not cached because it is used as the ODBC workaround for obtaining the last insert ID value. See the Connector/ODBC section of Chapter 27, *Connectors and APIs*.

  Statements within transactions that use `SERIALIZABLE` isolation level also cannot be cached because they use `LOCK IN SHARE MODE` locking.

- It uses `TEMPORARY` tables.
- It does not use any tables.
- It generates warnings.
- The user has a column-level privilege for any of the involved tables.

8.10.3.2 Query Cache SELECT Options

Two query cache-related options may be specified in `SELECT` statements:

- `SQL_CACHE`

  The query result is cached if it is cacheable and the value of the `query_cache_type` system variable is `ON` or `DEMAND`.

- `SQL_NO_CACHE`

  The server does not use the query cache. It neither checks the query cache to see whether the result is already cached, nor does it cache the query result.

Examples:

```
SELECT SQL_CACHE id, name FROM customer;
SELECT SQL_NO_CACHE id, name FROM customer;
```

8.10.3.3 Query Cache Configuration

The `have_query_cache` server system variable indicates whether the query cache is available:

```
mysql> SHOW VARIABLES LIKE 'have_query_cache';
+------------------+-------+
| Variable_name    | Value |
+------------------+-------+
| have_query_cache | YES   |
+------------------+-------+
```

When using a standard MySQL binary, this value is always `YES`, even if query caching is disabled.

Several other system variables control query cache operation. These can be set in an option file or on the command line when starting `mysqld`. The query cache system variables all have names that begin with `query_cache_`. They are described briefly in Section 5.1.5, “Server System Variables”, with additional configuration information given here.

To set the size of the query cache, set the `query_cache_size` system variable. Setting it to 0 disables the query cache, as does setting `query_cache_type=OFF`. By default, the query cache is disabled, achieved using defaults of `query_cache_size=0` and `query_cache_type=0`.

To reduce overhead significantly, start the server with `query_cache_type=0` if you will not be using the query cache.

Note

When using the Windows Configuration Wizard to install or configure MySQL, the default value for `query_cache_size` will be configured automatically for you based on the different configuration types available. When using the Windows Configuration Wizard, the query cache may be enabled (that is, set to a nonzero value) due to the selected configuration. The query cache is also controlled by the setting of the `query_cache_type` variable. Check the values of these variables as set in your `my.ini` file after configuration has taken place.

When you set `query_cache_size` to a nonzero value, keep in mind that the query cache needs a minimum size of about 40KB to allocate its structures. (The exact size depends on system architecture.) If you set the value too small, you'll get a warning, as in this example:

```
mysql> SET GLOBAL query_cache_size = 40000;
Query OK, 0 rows affected, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1282
Message: Query cache failed to set size 39936;
         new query cache size is 0

mysql> SET GLOBAL query_cache_size = 41984;
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW VARIABLES LIKE 'query_cache_size';
+------------------+-------+
| Variable_name    | Value |
+------------------+-------+
| query_cache_size | 41984 |
+------------------+-------+
```

For the query cache to actually be able to hold any query results, its size must be set larger:

```
mysql> SET GLOBAL query_cache_size = 1000000;
Query OK, 0 rows affected (0.04 sec)

mysql> SHOW VARIABLES LIKE 'query_cache_size';
+------------------+--------+
| Variable_name    | Value  |
+------------------+--------+
| query_cache_size | 999424 |
+------------------+--------+
1 row in set (0.00 sec)
```

The `query_cache_size` value is aligned to the nearest 1024 byte block. The value reported may therefore be different from the value that you assign.

If the query cache size is greater than 0, the `query_cache_type` variable influences how it works. This variable can be set to the following values:

- A value of `0` or `OFF` prevents caching or retrieval of cached results.
- A value of `1` or `ON` enables caching except of those statements that begin with `SELECT SQL_NO_CACHE`.
- A value of `2` or `DEMAND` causes caching of only those statements that begin with `SELECT SQL_CACHE`.

If `query_cache_size` is 0, you should also set `query_cache_type` variable to 0. In this case, the server does not acquire the query cache mutex at all, which means that the query cache cannot be enabled at runtime and there is reduced overhead in query execution.

Setting the `GLOBAL query_cache_type` value determines query cache behavior for all clients that connect after the change is made. Individual clients can control cache behavior for their own connection by setting the `SESSION query_cache_type` value. For example, a client can disable use of the query cache for its own queries like this:

```
mysql> SET SESSION query_cache_type = OFF;
```

If you set `query_cache_type` at server startup (rather than at runtime with a `SET` statement), only the numeric values are permitted.

To control the maximum size of individual query results that can be cached, set the `query_cache_limit` system variable. The default value is 1MB.

Be careful not to set the size of the cache too large. Due to the need for threads to lock the cache during updates, you may see lock contention issues with a very large cache.

Note

You can set the maximum size that can be specified for the query cache at runtime with the `SET` statement by using the `--maximum-query_cache_size=32M` option on the command line or in the configuration file.

When a query is to be cached, its result (the data sent to the client) is stored in the query cache during result retrieval. Therefore the data usually is not handled in one big chunk. The query cache allocates blocks for storing this data on demand, so when one block is filled, a new block is allocated. Because memory allocation operation is costly (timewise), the query cache allocates blocks with a minimum size given by the `query_cache_min_res_unit` system variable. When a query is executed, the last result block is trimmed to the actual data size so that unused memory is freed. Depending on the types of queries your server executes, you might find it helpful to tune the value of `query_cache_min_res_unit`:

- The default value of `query_cache_min_res_unit` is 4KB. This should be adequate for most cases.
- If you have a lot of queries with small results, the default block size may lead to memory fragmentation, as indicated by a large number of free blocks. Fragmentation can force the query cache to prune (delete) queries from the cache due to lack of memory. In this case, decrease the value of `query_cache_min_res_unit`. The number of free blocks and queries removed due to pruning are given by the values of the `Qcache_free_blocks` and `Qcache_lowmem_prunes` status variables.
- If most of your queries have large results (check the `Qcache_total_blocks` and `Qcache_queries_in_cache` status variables), you can increase performance by increasing `query_cache_min_res_unit`. However, be careful to not make it too large (see the previous item).

8.10.3.4 Query Cache Status and Maintenance

To check whether the query cache is present in your MySQL server, use the following statement:

```
mysql> SHOW VARIABLES LIKE 'have_query_cache';
+------------------+-------+
| Variable_name    | Value |
+------------------+-------+
| have_query_cache | YES   |
+------------------+-------+
```

You can defragment the query cache to better utilize its memory with the `FLUSH QUERY CACHE` statement. The statement does not remove any queries from the cache.

The `RESET QUERY CACHE` statement removes all query results from the query cache. The `FLUSH TABLES` statement also does this.

To monitor query cache performance, use `SHOW STATUS` to view the cache status variables:

```
mysql> SHOW STATUS LIKE 'Qcache%';
+-------------------------+--------+
| Variable_name           | Value  |
+-------------------------+--------+
| Qcache_free_blocks      | 36     |
```

```
| Qcache_free_memory      | 138488 |
| Qcache_hits             | 79570  |
| Qcache_inserts          | 27087  |
| Qcache_lowmem_prunes    | 3114   |
| Qcache_not_cached       | 22989  |
| Qcache_queries_in_cache | 415    |
| Qcache_total_blocks     | 912    |
+-------------------------+--------+
```

Descriptions of each of these variables are given in Section 5.1.7, “Server Status Variables”. Some uses for them are described here.

The total number of `SELECT` queries is given by this formula:

```
  Com_select
+ Qcache_hits
+ queries with errors found by parser
```

The `Com_select` value is given by this formula:

```
  Qcache_inserts
+ Qcache_not_cached
+ queries with errors found during the column-privileges check
```

The query cache uses variable-length blocks, so `Qcache_total_blocks` and `Qcache_free_blocks` may indicate query cache memory fragmentation. After `FLUSH QUERY CACHE`, only a single free block remains.

Every cached query requires a minimum of two blocks (one for the query text and one or more for the query results). Also, every table that is used by a query requires one block. However, if two or more queries use the same table, only one table block needs to be allocated.

The information provided by the `Qcache_lowmem_prunes` status variable can help you tune the query cache size. It counts the number of queries that have been removed from the cache to free up memory for caching new queries. The query cache uses a least recently used (LRU) strategy to decide which queries to remove from the cache. Tuning information is given in Section 8.10.3.3, “Query Cache Configuration”.

8.10.4 Caching of Prepared Statements and Stored Programs

For certain statements that a client might execute multiple times during a session, the server converts the statement to an internal structure and caches that structure to be used during execution. Caching enables the server to perform more efficiently because it avoids the overhead of reconverting the statement should it be needed again during the session. Conversion and caching occurs for these statements:

- Prepared statements, both those processed at the SQL level (using the `PREPARE` statement) and those processed using the binary client/server protocol (using the `mysql_stmt_prepare()` C API function). The `max_prepared_stmt_count` system variable controls the total number of statements the server caches. (The sum of the number of prepared statements across all sessions.)
- Stored programs (stored procedures and functions, triggers, and events). In this case, the server converts and caches the entire program body. The `stored_program_cache` system variable indicates the approximate number of stored programs the server caches per session.

The server maintains caches for prepared statements and stored programs on a per-session basis. Statements cached for one session are not accessible to other sessions. When a session ends, the server discards any statements cached for it.

When the server uses a cached internal statement structure, it must take care that the structure does not go out of date. Metadata changes can occur for an object used by the statement, causing

a mismatch between the current object definition and the definition as represented in the internal statement structure. Metadata changes occur for DDL statements such as those that create, drop, alter, rename, or truncate tables, or that analyze, optimize, or repair tables. Table content changes (for example, with `INSERT` or `UPDATE`) do not change metadata, nor do `SELECT` statements.

Here is an illustration of the problem. Suppose that a client prepares this statement:

```
PREPARE s1 FROM 'SELECT * FROM t1';
```

The `SELECT *` expands in the internal structure to the list of columns in the table. If the set of columns in the table is modified with `ALTER TABLE`, the prepared statement goes out of date. If the server does not detect this change the next time the client executes `s1`, the prepared statement will return incorrect results.

To avoid problems caused by metadata changes to tables or views referred to by the prepared statement, the server detects these changes and automatically reprepares the statement when it is next executed. That is, the server reparses the statement and rebuilds the internal structure. Reparsing also occurs after referenced tables or views are flushed from the table definition cache, either implicitly to make room for new entries in the cache, or explicitly due to `FLUSH TABLES`.

Similarly, if changes occur to objects used by a stored program, the server reparses affected statements within the program.

The server also detects metadata changes for objects in expressions. These might be used in statements specific to stored programs, such as `DECLARE CURSOR` or flow-control statements such as `IF`, `CASE`, and `RETURN`.

To avoid reparsing entire stored programs, the server reparses affected statements or expressions within a program only as needed. Examples:

- Suppose that metadata for a table or view is changed. Reparsing occurs for a `SELECT *` within the program that accesses the table or view, but not for a `SELECT *` that does not access the table or view.
- When a statement is affected, the server reparses it only partially if possible. Consider this `CASE` statement:

  ```
  CASE case_expr
    WHEN when_expr1 ...
    WHEN when_expr2 ...
    WHEN when_expr3 ...
    ...
  END CASE
  ```

  If a metadata change affects only `WHEN when_expr3`, that expression is reparsed. *`case_expr`* and the other `WHEN` expressions are not reparsed.

Reparsing uses the default database and SQL mode that were in effect for the original conversion to internal form.

The server attempts reparsing up to three times. An error occurs if all attempts fail.

Reparsing is automatic, but to the extent that it occurs, diminishes prepared statement and stored program performance.

For prepared statements, the `Com_stmt_reprepare` status variable tracks the number of repreparations.

8.11 Optimizing Locking Operations

MySQL manages contention for table contents using locking:

- Internal locking is performed within the MySQL server itself to manage contention for table contents by multiple threads. This type of locking is internal because it is performed entirely by the server and involves no other programs. See Section 8.11.1, “Internal Locking Methods”.
- External locking occurs when the server and other programs lock `MyISAM` table files to coordinate among themselves which program can access the tables at which time. See Section 8.11.5, “External Locking”.

8.11.1 Internal Locking Methods

This section discusses internal locking; that is, locking performed within the MySQL server itself to manage contention for table contents by multiple sessions. This type of locking is internal because it is performed entirely by the server and involves no other programs. For locking performed on MySQL files by other programs, see Section 8.11.5, “External Locking”.

- Row-Level Locking
- Table-Level Locking
- Choosing the Type of Locking

Row-Level Locking

MySQL uses row-level locking for `InnoDB` tables to support simultaneous write access by multiple sessions, making them suitable for multi-user, highly concurrent, and OLTP applications.

To avoid deadlocks when performing multiple concurrent write operations on a single `InnoDB` table, acquire necessary locks at the start of the transaction by issuing a `SELECT ... FOR UPDATE` statement for each group of rows expected to be modified, even if the data change statements come later in the transaction. If transactions modify or lock more than one table, issue the applicable statements in the same order within each transaction. Deadlocks affect performance rather than representing a serious error, because `InnoDB` automatically detects deadlock conditions by default and rolls back one of the affected transactions.

On high concurrency systems, deadlock detection can cause a slowdown when numerous threads wait for the same lock. At times, it may be more efficient to disable deadlock detection and rely on the `innodb_lock_wait_timeout` setting for transaction rollback when a deadlock occurs. Deadlock detection can be disabled using the `innodb_deadlock_detect` configuration option.

Advantages of row-level locking:

- Fewer lock conflicts when different sessions access different rows.
- Fewer changes for rollbacks.
- Possible to lock a single row for a long time.

Table-Level Locking

MySQL uses table-level locking for `MyISAM`, `MEMORY`, and `MERGE` tables, permitting only one session to update those tables at a time. This locking level makes these storage engines more suitable for read-only, read-mostly, or single-user applications.

These storage engines avoid deadlocks by always requesting all needed locks at once at the beginning of a query and always locking the tables in the same order. The tradeoff is that this strategy reduces concurrency; other sessions that want to modify the table must wait until the current data change statement finishes.

Advantages of table-level locking:

- Relatively little memory required (row locking requires memory per row or group of rows locked)
- Fast when used on a large part of the table because only a single lock is involved.
- Fast if you often do `GROUP BY` operations on a large part of the data or must scan the entire table frequently.

MySQL grants table write locks as follows:

1. If there are no locks on the table, put a write lock on it.
2. Otherwise, put the lock request in the write lock queue.

MySQL grants table read locks as follows:

1. If there are no write locks on the table, put a read lock on it.
2. Otherwise, put the lock request in the read lock queue.

Table updates are given higher priority than table retrievals. Therefore, when a lock is released, the lock is made available to the requests in the write lock queue and then to the requests in the read lock queue. This ensures that updates to a table are not "starved" even when there is heavy `SELECT` activity for the table. However, if there are many updates for a table, `SELECT` statements wait until there are no more updates.

For information on altering the priority of reads and writes, see Section 8.11.2, "Table Locking Issues".

You can analyze the table lock contention on your system by checking the `Table_locks_immediate` and `Table_locks_waited` status variables, which indicate the number of times that requests for table locks could be granted immediately and the number that had to wait, respectively:

```
mysql> SHOW STATUS LIKE 'Table%';
+-----------------------+---------+
| Variable_name         | Value   |
+-----------------------+---------+
| Table_locks_immediate | 1151552 |
| Table_locks_waited    | 15324   |
+-----------------------+---------+
```

The Performance Schema lock tables also provide locking information. See Section 25.11.12, "Performance Schema Lock Tables".

The `MyISAM` storage engine supports concurrent inserts to reduce contention between readers and writers for a given table: If a `MyISAM` table has no free blocks in the middle of the data file, rows are always inserted at the end of the data file. In this case, you can freely mix concurrent `INSERT` and `SELECT` statements for a `MyISAM` table without locks. That is, you can insert rows into a `MyISAM` table at the same time other clients are reading from it. Holes can result from rows having been deleted from or updated in the middle of the table. If there are holes, concurrent inserts are disabled but are enabled again automatically when all holes have been filled with new data. To control this behavior, use the `concurrent_insert` system variable. See Section 8.11.3, "Concurrent Inserts".

If you acquire a table lock explicitly with `LOCK TABLES`, you can request a `READ LOCAL` lock rather than a `READ` lock to enable other sessions to perform concurrent inserts while you have the table locked.

To perform many `INSERT` and `SELECT` operations on a table `t1` when concurrent inserts are not possible, you can insert rows into a temporary table `temp_t1` and update the real table with the rows from the temporary table:

```
mysql> LOCK TABLES t1 WRITE, temp_t1 WRITE;
mysql> INSERT INTO t1 SELECT * FROM temp_t1;
```

```
mysql> DELETE FROM temp_t1;
mysql> UNLOCK TABLES;
```

Choosing the Type of Locking

Generally, table locks are superior to row-level locks in the following cases:

- Most statements for the table are reads.
- Statements for the table are a mix of reads and writes, where writes are updates or deletes for a single row that can be fetched with one key read:

```
UPDATE tbl_name SET column=value WHERE unique_key_col=key_value;
DELETE FROM tbl_name WHERE unique_key_col=key_value;
```

- `SELECT` combined with concurrent `INSERT` statements, and very few `UPDATE` or `DELETE` statements.
- Many scans or `GROUP BY` operations on the entire table without any writers.

With higher-level locks, you can more easily tune applications by supporting locks of different types, because the lock overhead is less than for row-level locks.

Options other than row-level locking:

- Versioning (such as that used in MySQL for concurrent inserts) where it is possible to have one writer at the same time as many readers. This means that the database or table supports different views for the data depending on when access begins. Other common terms for this are “time travel,” “copy on write,” or “copy on demand.”
- Copy on demand is in many cases superior to row-level locking. However, in the worst case, it can use much more memory than using normal locks.
- Instead of using row-level locks, you can employ application-level locks, such as those provided by `GET_LOCK()` and `RELEASE_LOCK()` in MySQL. These are advisory locks, so they work only with applications that cooperate with each other. See Section 12.20, “Miscellaneous Functions”.

8.11.2 Table Locking Issues

`InnoDB` tables use row-level locking so that multiple sessions and applications can read from and write to the same table simultaneously, without making each other wait or producing inconsistent results. For this storage engine, avoid using the `LOCK TABLES` statement, because it does not offer any extra protection, but instead reduces concurrency. The automatic row-level locking makes these tables suitable for your busiest databases with your most important data, while also simplifying application logic since you do not need to lock and unlock tables. Consequently, the `InnoDB` storage engine is the default in MySQL.

MySQL uses table locking (instead of page, row, or column locking) for all storage engines except `InnoDB`. The locking operations themselves do not have much overhead. But because only one session can write to a table at any one time, for best performance with these other storage engines, use them primarily for tables that are queried often and rarely inserted into or updated.

- Performance Considerations Favoring InnoDB
- Workarounds for Locking Performance Issues

Performance Considerations Favoring InnoDB

When choosing whether to create a table using `InnoDB` or a different storage engine, keep in mind the following disadvantages of table locking:

- Table locking enables many sessions to read from a table at the same time, but if a session wants to write to a table, it must first get exclusive access, meaning it might have to wait for other sessions to finish with the table first. During the update, all other sessions that want to access this particular table must wait until the update is done.
- Table locking causes problems when a session is waiting because the disk is full and free space needs to become available before the session can proceed. In this case, all sessions that want to access the problem table are also put in a waiting state until more disk space is made available.
- A `SELECT` statement that takes a long time to run prevents other sessions from updating the table in the meantime, making the other sessions appear slow or unresponsive. While a session is waiting to get exclusive access to the table for updates, other sessions that issue `SELECT` statements will queue up behind it, reducing concurrency even for read-only sessions.

Workarounds for Locking Performance Issues

The following items describe some ways to avoid or reduce contention caused by table locking:

- Consider switching the table to the `InnoDB` storage engine, either using `CREATE TABLE ... ENGINE=INNODB` during setup, or using `ALTER TABLE ... ENGINE=INNODB` for an existing table. See Chapter 15, *The InnoDB Storage Engine* for more details about this storage engine.
- Optimize `SELECT` statements to run faster so that they lock tables for a shorter time. You might have to create some summary tables to do this.
- Start `mysqld` with `--low-priority-updates`. For storage engines that use only table-level locking (such as `MyISAM`, `MEMORY`, and `MERGE`), this gives all statements that update (modify) a table lower priority than `SELECT` statements. In this case, the second `SELECT` statement in the preceding scenario would execute before the `UPDATE` statement, and would not wait for the first `SELECT` to finish.
- To specify that all updates issued in a specific connection should be done with low priority, set the `low_priority_updates` server system variable equal to 1.
- To give a specific `INSERT`, `UPDATE`, or `DELETE` statement lower priority, use the `LOW_PRIORITY` attribute.
- To give a specific `SELECT` statement higher priority, use the `HIGH_PRIORITY` attribute. See Section 13.2.10, “SELECT Syntax”.
- Start `mysqld` with a low value for the `max_write_lock_count` system variable to force MySQL to temporarily elevate the priority of all `SELECT` statements that are waiting for a table after a specific number of inserts to the table occur. This permits `READ` locks after a certain number of `WRITE` locks.
- If you have problems with `INSERT` combined with `SELECT`, consider switching to `MyISAM` tables, which support concurrent `SELECT` and `INSERT` statements. (See Section 8.11.3, “Concurrent Inserts”.)
- If you have problems with mixed `SELECT` and `DELETE` statements, the `LIMIT` option to `DELETE` may help. See Section 13.2.2, “DELETE Syntax”.
- Using `SQL_BUFFER_RESULT` with `SELECT` statements can help to make the duration of table locks shorter. See Section 13.2.10, “SELECT Syntax”.
- Splitting table contents into separate tables may help, by allowing queries to run against columns in one table, while updates are confined to columns in a different table.
- You could change the locking code in `mysys/thr_lock.c` to use a single queue. In this case, write locks and read locks would have the same priority, which might help some applications.

8.11.3 Concurrent Inserts

The MyISAM storage engine supports concurrent inserts to reduce contention between readers and writers for a given table: If a MyISAM table has no holes in the data file (deleted rows in the middle), an INSERT statement can be executed to add rows to the end of the table at the same time that SELECT statements are reading rows from the table. If there are multiple INSERT statements, they are queued and performed in sequence, concurrently with the SELECT statements. The results of a concurrent INSERT may not be visible immediately.

The concurrent_insert system variable can be set to modify the concurrent-insert processing. By default, the variable is set to AUTO (or 1) and concurrent inserts are handled as just described. If concurrent_insert is set to NEVER (or 0), concurrent inserts are disabled. If the variable is set to ALWAYS (or 2), concurrent inserts at the end of the table are permitted even for tables that have deleted rows. See also the description of the concurrent_insert system variable.

If you are using the binary log, concurrent inserts are converted to normal inserts for CREATE ... SELECT or INSERT ... SELECT statements. This is done to ensure that you can re-create an exact copy of your tables by applying the log during a backup operation. See Section 5.4.4, "The Binary Log". In addition, for those statements a read lock is placed on the selected-from table such that inserts into that table are blocked. The effect is that concurrent inserts for that table must wait as well.

With LOAD DATA INFILE, if you specify CONCURRENT with a MyISAM table that satisfies the condition for concurrent inserts (that is, it contains no free blocks in the middle), other sessions can retrieve data from the table while LOAD DATA is executing. Use of the CONCURRENT option affects the performance of LOAD DATA a bit, even if no other session is using the table at the same time.

If you specify HIGH_PRIORITY, it overrides the effect of the --low-priority-updates option if the server was started with that option. It also causes concurrent inserts not to be used.

For LOCK TABLE, the difference between READ LOCAL and READ is that READ LOCAL permits nonconflicting INSERT statements (concurrent inserts) to execute while the lock is held. However, this cannot be used if you are going to manipulate the database using processes external to the server while you hold the lock.

8.11.4 Metadata Locking

MySQL uses metadata locking to manage concurrent access to database objects and to ensure data consistency. Metadata locking applies not just to tables, but also to schemas, stored programs (procedures, functions, triggers, and scheduled events), and tablespaces.

Metadata locking does involve some overhead, which increases as query volume increases. Metadata contention increases the more that multiple queries attempt to access the same objects.

Metadata locking is not a replacement for the table definition cache, and its mutexes and locks differ from the LOCK_open mutex. The following discussion provides some information about how metadata locking works.

To ensure transaction serializability, the server must not permit one session to perform a data definition language (DDL) statement on a table that is used in an uncompleted explicitly or implicitly started transaction in another session. The server achieves this by acquiring metadata locks on tables used within a transaction and deferring release of those locks until the transaction ends. A metadata lock on a table prevents changes to the table's structure. This locking approach has the implication that a table that is being used by a transaction within one session cannot be used in DDL statements by other sessions until the transaction ends.

This principle applies not only to transactional tables, but also to nontransactional tables. Suppose that a session begins a transaction that uses transactional table t and nontransactional table nt as follows:

```
START TRANSACTION;
SELECT * FROM t;
SELECT * FROM nt;
```

The server holds metadata locks on both `t` and `nt` until the transaction ends. If another session attempts a DDL or write lock operation on either table, it blocks until metadata lock release at transaction end. For example, a second session blocks if it attempts any of these operations:

```
DROP TABLE t;
ALTER TABLE t ...;
DROP TABLE nt;
ALTER TABLE nt ...;
LOCK TABLE t ... WRITE;
```

The same behavior applies for The `LOCK TABLES ... READ`. That is, explicitly or implicitly started transactions that update any table (transactional or nontransactional) will block and be blocked by `LOCK TABLES ... READ` for that table.

If the server acquires metadata locks for a statement that is syntactically valid but fails during execution, it does not release the locks early. Lock release is still deferred to the end of the transaction because the failed statement is written to the binary log and the locks protect log consistency.

In autocommit mode, each statement is in effect a complete transaction, so metadata locks acquired for the statement are held only to the end of the statement.

Metadata locks acquired during a `PREPARE` statement are released once the statement has been prepared, even if preparation occurs within a multiple-statement transaction.

8.11.5 External Locking

External locking is the use of file system locking to manage contention for `MyISAM` database tables by multiple processes. External locking is used in situations where a single process such as the MySQL server cannot be assumed to be the only process that requires access to tables. Here are some examples:

- If you run multiple servers that use the same database directory (not recommended), each server must have external locking enabled.
- If you use `myisamchk` to perform table maintenance operations on `MyISAM` tables, you must either ensure that the server is not running, or that the server has external locking enabled so that it locks table files as necessary to coordinate with `myisamchk` for access to the tables. The same is true for use of `myisampack` to pack `MyISAM` tables.

  If the server is run with external locking enabled, you can use `myisamchk` at any time for read operations such a checking tables. In this case, if the server tries to update a table that `myisamchk` is using, the server will wait for `myisamchk` to finish before it continues.

  If you use `myisamchk` for write operations such as repairing or optimizing tables, or if you use `myisampack` to pack tables, you *must* always ensure that the `mysqld` server is not using the table. If you do not stop `mysqld`, at least do a `mysqladmin flush-tables` before you run `myisamchk`. Your tables *may become corrupted* if the server and `myisamchk` access the tables simultaneously.

With external locking in effect, each process that requires access to a table acquires a file system lock for the table files before proceeding to access the table. If all necessary locks cannot be acquired, the process is blocked from accessing the table until the locks can be obtained (after the process that currently holds the locks releases them).

External locking affects server performance because the server must sometimes wait for other processes before it can access tables.

External locking is unnecessary if you run a single server to access a given data directory (which is the usual case) and if no other programs such as `myisamchk` need to modify tables while the server is running. If you only *read* tables with other programs, external locking is not required, although `myisamchk` might report warnings if the server changes tables while `myisamchk` is reading them.

With external locking disabled, to use `myisamchk`, you must either stop the server while `myisamchk` executes or else lock and flush the tables before running `myisamchk`. (See System Factors.) To avoid this requirement, use the `CHECK TABLE` and `REPAIR TABLE` statements to check and repair `MyISAM` tables.

For `mysqld`, external locking is controlled by the value of the `skip_external_locking` system variable. When this variable is enabled, external locking is disabled, and vice versa. External locking is disabled by default.

Use of external locking can be controlled at server startup by using the `--external-locking` or `--skip-external-locking` option.

If you do use external locking option to enable updates to `MyISAM` tables from many MySQL processes, you must ensure that the following conditions are satisfied:

- Do not use the query cache for queries that use tables that are updated by another process.
- Do not start the server with the `--delay-key-write=ALL` option or use the `DELAY_KEY_WRITE=1` table option for any shared tables. Otherwise, index corruption can occur.

The easiest way to satisfy these conditions is to always use `--external-locking` together with `--delay-key-write=OFF` and `--query-cache-size=0`. (This is not done by default because in many setups it is useful to have a mixture of the preceding options.)

8.12 Optimizing the MySQL Server

This section discusses optimization techniques for the database server, primarily dealing with system configuration rather than tuning SQL statements. The information in this section is appropriate for DBAs who want to ensure performance and scalability across the servers they manage; for developers constructing installation scripts that include setting up the database; and people running MySQL themselves for development, testing, and so on who want to maximize their own productivity.

8.12.1 Optimizing Disk I/O

This section describes ways to configure storage devices when you can devote more and faster storage hardware to the database server. For information about optimizing an `InnoDB` configuration to improve I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- Disk seeks are a huge performance bottleneck. This problem becomes more apparent when the amount of data starts to grow so large that effective caching becomes impossible. For large databases where you access data more or less randomly, you can be sure that you need at least one disk seek to read and a couple of disk seeks to write things. To minimize this problem, use disks with low seek times.
- Increase the number of available disk spindles (and thereby reduce the seek overhead) by either symlinking files to different disks or striping the disks:
  - Using symbolic links

    This means that, for `MyISAM` tables, you symlink the index file and data files from their usual location in the data directory to another disk (that may also be striped). This makes both the seek and read times better, assuming that the disk is not used for other purposes as well. See Section 8.12.2, “Using Symbolic Links”.

    Symbolic links are not supported for use with `InnoDB` tables. However, it is possible to place `InnoDB` data and log files on different physical disks. For more information, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

  - Striping

Striping means that you have many disks and put the first block on the first disk, the second block on the second disk, and the *`N`*-th block on the (*`N MOD number_of_disks`*) disk, and so on. This means if your normal data size is less than the stripe size (or perfectly aligned), you get much better performance. Striping is very dependent on the operating system and the stripe size, so benchmark your application with different stripe sizes. See Section 8.13.2, “Using Your Own Benchmarks”.

The speed difference for striping is *very* dependent on the parameters. Depending on how you set the striping parameters and number of disks, you may get differences measured in orders of magnitude. You have to choose to optimize for random or sequential access.

- For reliability, you may want to use RAID 0+1 (striping plus mirroring), but in this case, you need 2 × *`N`* drives to hold *`N`* drives of data. This is probably the best option if you have the money for it. However, you may also have to invest in some volume-management software to handle it efficiently.
- A good option is to vary the RAID level according to how critical a type of data is. For example, store semi-important data that can be regenerated on a RAID 0 disk, but store really important data such as host information and logs on a RAID 0+1 or RAID *`N`* disk. RAID *`N`* can be a problem if you have many writes, due to the time required to update the parity bits.
- You can also set the parameters for the file system that the database uses:

  If you do not need to know when files were last accessed (which is not really useful on a database server), you can mount your file systems with the `-o noatime` option. That skips updates to the last access time in inodes on the file system, which avoids some disk seeks.

  On many operating systems, you can set a file system to be updated asynchronously by mounting it with the `-o async` option. If your computer is reasonably stable, this should give you better performance without sacrificing too much reliability. (This flag is on by default on Linux.)

Using NFS with MySQL

Caution is advised when considering using NFS with MySQL. Potential issues, which vary by operating system and NFS version, include:

- MySQL data and log files placed on NFS volumes becoming locked and unavailable for use. Locking issues may occur in cases where multiple instances of MySQL access the same data directory or where MySQL is shut down improperly, due to a power outage, for example. NFS version 4 addresses underlying locking issues with the introduction of advisory and lease-based locking. However, sharing a data directory among MySQL instances is not recommended.
- Data inconsistencies introduced due to messages received out of order or lost network traffic. To avoid this issue, use TCP with `hard` and `intr` mount options.
- Maximum file size limitations. NFS Version 2 clients can only access the lowest 2GB of a file (signed 32 bit offset). NFS Version 3 clients support larger files (up to 64 bit offsets). The maximum supported file size also depends on the local file system of the NFS server.

Using NFS within a professional SAN environment or other storage system tends to offer greater reliability than using NFS outside of such an environment. However, NFS within a SAN environment may be slower than directly attached or bus-attached non-rotational storage.

If you choose to use NFS, NFS Version 4 or later is recommended, as is testing your NFS setup thoroughly before deploying into a production environment.

8.12.2 Using Symbolic Links

You can move databases or tables from the database directory to other locations and replace them with symbolic links to the new locations. You might want to do this, for example, to move a database to a file system with more free space or increase the speed of your system by spreading your tables to different disks.

For InnoDB tables, use the DATA DIRECTORY clause on the CREATE TABLE statement instead of symbolic links, as explained in Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory”. This new feature is a supported, cross-platform technique.

The recommended way to do this is to symlink entire database directories to a different disk. Symlink MyISAM tables only as a last resort.

To determine the location of your data directory, use this statement:

```
SHOW VARIABLES LIKE 'datadir';
```

8.12.2.1 Using Symbolic Links for Databases on Unix

On Unix, the way to symlink a database is first to create a directory on some disk where you have free space and then to create a soft link to it from the MySQL data directory.

```
shell> mkdir /dr1/databases/test
shell> ln -s /dr1/databases/test /path/to/datadir
```

MySQL does not support linking one directory to multiple databases. Replacing a database directory with a symbolic link works as long as you do not make a symbolic link between databases. Suppose that you have a database db1 under the MySQL data directory, and then make a symlink db2 that points to db1:

```
shell> cd /path/to/datadir
shell> ln -s db1 db2
```

The result is that, for any table tbl_a in db1, there also appears to be a table tbl_a in db2. If one client updates db1.tbl_a and another client updates db2.tbl_a, problems are likely to occur.

8.12.2.2 Using Symbolic Links for MyISAM Tables on Unix

Note

As of MySQL 8.0.2, symbolic link support as described here, along with the the --symbolic-links option that controls it, is deprecated and will be removed in a future version of MySQL. In addition, the option is disabled by default.

Symlinks are fully supported only for MyISAM tables. For files used by tables for other storage engines, you may get strange problems if you try to use symbolic links. For InnoDB tables, use the alternative technique explained in Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory” instead.

Do not symlink tables on systems that do not have a fully operational realpath() call. (Linux and Solaris support realpath()). To determine whether your system supports symbolic links, check the value of the have_symlink system variable using this statement:

```
SHOW VARIABLES LIKE 'have_symlink';
```

The handling of symbolic links for MyISAM tables works as follows:

- In the data directory, you always have the data (.MYD) file and the index (.MYI) file. The data file and index file can be moved elsewhere and replaced in the data directory by symlinks.
- You can symlink the data file and the index file independently to different directories.
- To instruct a running MySQL server to perform the symlinking, use the DATA DIRECTORY and INDEX DIRECTORY options to CREATE TABLE. See Section 13.1.15, “CREATE TABLE Syntax”. Alternatively, if mysqld is not running, symlinking can be accomplished manually using ln -s from the command line.

Note

The path used with either or both of the `DATA DIRECTORY` and `INDEX DIRECTORY` options may not include the MySQL `data` directory. (Bug #32167)

- `myisamchk` does not replace a symlink with the data file or index file. It works directly on the file to which the symlink points. Any temporary files are created in the directory where the data file or index file is located. The same is true for the `ALTER TABLE`, `OPTIMIZE TABLE`, and `REPAIR TABLE` statements.

-

  Note

  When you drop a table that is using symlinks, *both the symlink and the file to which the symlink points are dropped*. This is an extremely good reason *not* to run `mysqld` as the system `root` or permit system users to have write access to MySQL database directories.

- If you rename a table with `ALTER TABLE ... RENAME` or `RENAME TABLE` and you do not move the table to another database, the symlinks in the database directory are renamed to the new names and the data file and index file are renamed accordingly.

- If you use `ALTER TABLE ... RENAME` or `RENAME TABLE` to move a table to another database, the table is moved to the other database directory. If the table name changed, the symlinks in the new database directory are renamed to the new names and the data file and index file are renamed accordingly.

- If you are not using symlinks, start `mysqld` with the `--skip-symbolic-links` option to ensure that no one can use `mysqld` to drop or rename a file outside of the data directory.

These table symlink operations are not supported:

- `ALTER TABLE` ignores the `DATA DIRECTORY` and `INDEX DIRECTORY` table options.

8.12.2.3 Using Symbolic Links for Databases on Windows

On Windows, symbolic links can be used for database directories. This enables you to put a database directory at a different location (for example, on a different disk) by setting up a symbolic link to it. Use of database symlinks on Windows is similar to their use on Unix, although the procedure for setting up the link differs.

Suppose that you want to place the database directory for a database named `mydb` at `D:\data\mydb`. To do this, create a symbolic link in the MySQL data directory that points to `D:\data\mydb`. However, before creating the symbolic link, make sure that the `D:\data\mydb` directory exists by creating it if necessary. If you already have a database directory named `mydb` in the data directory, move it to `D:\data`. Otherwise, the symbolic link will be ineffective. To avoid problems, make sure that the server is not running when you move the database directory.

Windows Vista, Windows Server 2008, or newer have native symbolic link support, so you can create a symlink using the `mklink` command. This command requires administrative privileges.

1. Change location into the data directory:

   ```
   C:\> cd \path\to\datadir
   ```

2. In the data directory, create a symlink named `mydb` that points to the location of the database directory:

   ```
   C:\> mklink /d mydb D:\data\mydb
   ```

After this, all tables created in the database `mydb` are created in `D:\data\mydb`.

8.12.3 Optimizing Memory Use

8.12.3.1 How MySQL Uses Memory

MySQL allocates buffers and caches to improve performance of database operations. The default configuration is designed to permit a MySQL server to start on a virtual machine that has approximately 512MB of RAM. You can improve MySQL performance by increasing the values of certain cache and buffer-related system variables. You can also modify the default configuration to run MySQL on systems with limited memory.

The following list describes some of the ways that MySQL uses memory. Where applicable, relevant system variables are referenced. Some items are storage engine or feature specific.

- The `InnoDB` buffer pool is a memory area that holds cached `InnoDB` data for tables, indexes, and other auxiliary buffers. For efficiency of high-volume read operations, the buffer pool is divided into pages that can potentially hold multiple rows. For efficiency of cache management, the buffer pool is implemented as a linked list of pages; data that is rarely used is aged out of the cache, using a variation of the LRU algorithm. For more information, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

  The size of the buffer pool is important for system performance:

  - `InnoDB` allocates memory for the entire buffer pool at server startup, using `malloc()` operations. The `innodb_buffer_pool_size` system variable defines the buffer pool size. Typically, a recommended `innodb_buffer_pool_size` value is 50 to 75 percent of system memory. `innodb_buffer_pool_size` can be configured dynamically, while the server is running. For more information, see Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”.

  - On systems with a large amount of memory, you can improve concurrency by dividing the buffer pool into multiple buffer pool instances. The `innodb_buffer_pool_instances` system variable defines the number of buffer pool instances.

  - A buffer pool that is too small may cause excessive churning as pages are flushed from the buffer pool only to be required again a short time later.

  - A buffer pool that is too large may cause swapping due to competition for memory.

- The storage engine interface enables the optimizer to provide information about the size of the record buffer to be used for scans that the optimizer estimates will read multiple rows. The buffer size can vary based on the size of the estimate. `InnoDB` uses this variable-size buffering capability to take advantage of row prefetching, and to reduce the overhead of latching and B-tree navigation.

- All threads share the `MyISAM` key buffer. The `key_buffer_size` system variable determines its size.

  For each `MyISAM` table the server opens, the index file is opened once; the data file is opened once for each concurrently running thread that accesses the table. For each concurrent thread, a table structure, column structures for each column, and a buffer of size `3 * N` are allocated (where *`N`* is the maximum row length, not counting `BLOB` columns). A `BLOB` column requires five to eight bytes plus the length of the `BLOB` data. The `MyISAM` storage engine maintains one extra row buffer for internal use.

- The `myisam_use_mmap` system variable can be set to 1 to enable memory-mapping for all `MyISAM` tables.

- If an internal in-memory temporary table becomes too large (as determined using the `tmp_table_size` and `max_heap_table_size` system variables), MySQL automatically converts the table from in-memory to on-disk format. On-disk temporary tables use the storage engine defined

by the `internal_tmp_disk_storage_engine` system variable. You can increase the permissible temporary table size as described in Section 8.4.4, “Internal Temporary Table Use in MySQL”.

For `MEMORY` tables explicitly created with `CREATE TABLE`, only the `max_heap_table_size` system variable determines how large the table is permitted to grow and there is no conversion to on-disk format.

- The MySQL Performance Schema is a feature for monitoring MySQL server execution at a low level. The Performance Schema dynamically allocates memory incrementally, scaling its memory use to actual server load, instead of allocating required memory during server startup. Once memory is allocated, it is not freed until the server is restarted. For more information, see Section 25.16, “The Performance Schema Memory-Allocation Model”.
- Each thread that the server uses to manage client connections requires some thread-specific space. The following list indicates these and which system variables control their size:
  - A stack (`thread_stack`)
  - A connection buffer (`net_buffer_length`)
  - A result buffer (`net_buffer_length`)

  The connection buffer and result buffer each begin with a size equal to `net_buffer_length` bytes, but are dynamically enlarged up to `max_allowed_packet` bytes as needed. The result buffer shrinks to `net_buffer_length` bytes after each SQL statement. While a statement is running, a copy of the current statement string is also allocated.

  Each connection thread uses memory for computing statement digests. The server allocates `max_digest_length` bytes per session. See Section 25.9, “Performance Schema Statement Digests”.
- All threads share the same base memory.
- When a thread is no longer needed, the memory allocated to it is released and returned to the system unless the thread goes back into the thread cache. In that case, the memory remains allocated.
- Each request that performs a sequential scan of a table allocates a *read buffer*. The `read_buffer_size` system variable determines the buffer size.
- When reading rows in an arbitrary sequence (for example, following a sort), a *random-read buffer* may be allocated to avoid disk seeks. The `read_rnd_buffer_size` system variable determines the buffer size.
- All joins are executed in a single pass, and most joins can be done without even using a temporary table. Most temporary tables are memory-based hash tables. Temporary tables with a large row length (calculated as the sum of all column lengths) or that contain `BLOB` columns are stored on disk.
- Most requests that perform a sort allocate a sort buffer and zero to two temporary files depending on the result set size. See Section B.5.3.5, “Where MySQL Stores Temporary Files”.
- Almost all parsing and calculating is done in thread-local and reusable memory pools. No memory overhead is needed for small items, thus avoiding the normal slow memory allocation and freeing. Memory is allocated only for unexpectedly large strings.
- For each table having `BLOB` columns, a buffer is enlarged dynamically to read in larger `BLOB` values. If you scan a table, the buffer grows as large as the largest `BLOB` value.
- MySQL requires memory and descriptors for the table cache. Handler structures for all in-use tables are saved in the table cache and managed as “First In, First Out” (FIFO). The `table_open_cache` system variable defines the initial table cache size; see Section 8.4.3.1, “How MySQL Opens and Closes Tables”.

MySQL also requires memory for the table definition cache. The `table_definition_cache` system variable defines the number of table definitions that can be stored in the table definition cache. If you use a large number of tables, you can create a large table definition cache to speed up the opening of tables. The table definition cache takes less space and does not use file descriptors, unlike the table cache.

- A `FLUSH TABLES` statement or `mysqladmin flush-tables` command closes all tables that are not in use at once and marks all in-use tables to be closed when the currently executing thread finishes. This effectively frees most in-use memory. `FLUSH TABLES` does not return until all tables have been closed.
- The server caches information in memory as a result of `GRANT`, `CREATE USER`, `CREATE SERVER`, and `INSTALL PLUGIN` statements. This memory is not released by the corresponding `REVOKE`, `DROP USER`, `DROP SERVER`, and `UNINSTALL PLUGIN` statements, so for a server that executes many instances of the statements that cause caching, there will be an increase in memory use. This cached memory can be freed with `FLUSH PRIVILEGES`.

`ps` and other system status programs may report that `mysqld` uses a lot of memory. This may be caused by thread stacks on different memory addresses. For example, the Solaris version of `ps` counts the unused memory between stacks as used memory. To verify this, check available swap with `swap -s`. We test `mysqld` with several memory-leakage detectors (both commercial and Open Source), so there should be no memory leaks.

Monitoring MySQL Memory Usage

The following example demonstrates how to use Performance Schema and sys schema to monitor MySQL memory usage.

Most Performance Schema memory instrumentation is disabled by default. Instruments can be enabled by updating the `ENABLED` column of the Performance Schema `setup_instruments` table. Memory instruments have names in the form of `memory/code_area/instrument_name`, where *`code_area`* is a value such as `sql` or `innodb`, and *`instrument_name`* is the instrument detail.

1. To view available MySQL memory instruments, query the Performance Schema `setup_instruments` table. The following query returns hundreds of memory instruments for all code areas.

```
mysql> SELECT * FROM performance_schema.setup_instruments
       WHERE NAME LIKE '%memory%';
```

You can narrow results by specifying a code area. For example, you can limit results to `InnoDB` memory instruments by specifying `innodb` as the code area.

```
mysql> SELECT * FROM performance_schema.setup_instruments
       WHERE NAME LIKE '%memory/innodb%';
+----------------------------------------+---------+-------+
| NAME                                   | ENABLED | TIMED |
+----------------------------------------+---------+-------+
| memory/innodb/adaptive hash index      | NO      | NO    |
| memory/innodb/buf_buf_pool             | NO      | NO    |
| memory/innodb/dict_stats_bg_recalc_pool_t | NO      | NO    |
| memory/innodb/dict_stats_index_map_t   | NO      | NO    |
| memory/innodb/dict_stats_n_diff_on_level | NO      | NO    |
| memory/innodb/other                    | NO      | NO    |
| memory/innodb/row_log_buf              | NO      | NO    |
| memory/innodb/row_merge_sort           | NO      | NO    |
| memory/innodb/std                      | NO      | NO    |
| memory/innodb/trx_sys_t::rw_trx_ids    | NO      | NO    |
...
```

Depending on your MySQL installation, code areas may include `performance_schema`, `sql`, `client`, `innodb`, `myisam`, `csv`, `memory`, `blackhole`, `archive`, `partition`, and others.

2. To enable memory instruments, add a `performance-schema-instrument` rule to your MySQL configuration file. For example, to enable all memory instruments, add this rule to your configuration file and restart the server:

```
performance-schema-instrument='memory/%=COUNTED'
```

Note

Enabling memory instruments at startup ensures that memory allocations that occur at startup are counted.

After restarting the server, the `ENABLED` column of the Performance Schema `setup_instruments` table should report `YES` for memory instruments that you enabled. The `TIMED` column in the `setup_instruments` table is ignored for memory instruments because memory operations are not timed.

```
mysql> SELECT * FROM performance_schema.setup_instruments
       WHERE NAME LIKE '%memory/innodb%';
+----------------------------------------+---------+-------+
| NAME                                   | ENABLED | TIMED |
+----------------------------------------+---------+-------+
| memory/innodb/adaptive hash index      | NO      | NO    |
| memory/innodb/buf_buf_pool             | NO      | NO    |
| memory/innodb/dict_stats_bg_recalc_pool_t | NO      | NO    |
| memory/innodb/dict_stats_index_map_t   | NO      | NO    |
| memory/innodb/dict_stats_n_diff_on_level | NO      | NO    |
| memory/innodb/other                    | NO      | NO    |
| memory/innodb/row_log_buf              | NO      | NO    |
| memory/innodb/row_merge_sort           | NO      | NO    |
| memory/innodb/std                      | NO      | NO    |
| memory/innodb/trx_sys_t::rw_trx_ids    | NO      | NO    |
...
```

3. Query memory instrument data. In this example, memory instrument data is queried in the Performance Schema `memory_summary_global_by_event_name` table, which summarizes data by `EVENT_NAME`. The `EVENT_NAME` is the name of the instrument.

   The following query returns memory data for the `InnoDB` buffer pool. For column descriptions, see Section 25.11.15.10, “Memory Summary Tables”.

```
mysql> SELECT * FROM performance_schema.memory_summary_global_by_event_name
       WHERE EVENT_NAME LIKE 'memory/innodb/buf_buf_pool'\G
                 EVENT_NAME: memory/innodb/buf_buf_pool
                COUNT_ALLOC: 1
                 COUNT_FREE: 0
   SUM_NUMBER_OF_BYTES_ALLOC: 137428992
    SUM_NUMBER_OF_BYTES_FREE: 0
              LOW_COUNT_USED: 0
          CURRENT_COUNT_USED: 1
             HIGH_COUNT_USED: 1
    LOW_NUMBER_OF_BYTES_USED: 0
CURRENT_NUMBER_OF_BYTES_USED: 137428992
   HIGH_NUMBER_OF_BYTES_USED: 137428992
```

   The same underlying data can be queried using the `sys` schema `memory_global_by_current_bytes` table, which shows current memory usage within the server globally, broken down by allocation type.

```
mysql> SELECT * FROM sys.memory_global_by_current_bytes
       WHERE event_name LIKE 'memory/innodb/buf_buf_pool'\G
*************************** 1. row ***************************
       event_name: memory/innodb/buf_buf_pool
    current_count: 1
    current_alloc: 131.06 MiB
current_avg_alloc: 131.06 MiB
       high_count: 1
       high_alloc: 131.06 MiB
```

```
   high_avg_alloc: 131.06 MiB
```

This `sys` schema query aggregates currently allocated memory (`current_alloc`) by code area:

```
mysql> SELECT SUBSTRING_INDEX(event_name,'/',2) AS
       code_area, sys.format_bytes(SUM(current_alloc))
       AS current_alloc
       FROM sys.x$memory_global_by_current_bytes
       GROUP BY SUBSTRING_INDEX(event_name,'/',2)
       ORDER BY SUM(current_alloc) DESC;
+---------------------------+---------------+
| code_area                 | current_alloc |
+---------------------------+---------------+
| memory/innodb             | 843.24 MiB    |
| memory/performance_schema | 81.29 MiB     |
| memory/mysys              | 8.20 MiB      |
| memory/sql                | 2.47 MiB      |
| memory/memory             | 174.01 KiB    |
| memory/myisam             | 46.53 KiB     |
| memory/blackhole          | 512 bytes     |
| memory/federated          | 512 bytes     |
| memory/csv                | 512 bytes     |
| memory/vio                | 496 bytes     |
+---------------------------+---------------+
```

For more information about `sys` schema, see Chapter 26, *MySQL sys Schema*.

8.12.3.2 Enabling Large Page Support

Some hardware/operating system architectures support memory pages greater than the default (usually 4KB). The actual implementation of this support depends on the underlying hardware and operating system. Applications that perform a lot of memory accesses may obtain performance improvements by using large pages due to reduced Translation Lookaside Buffer (TLB) misses.

In MySQL, large pages can be used by InnoDB, to allocate memory for its buffer pool and additional memory pool.

Standard use of large pages in MySQL attempts to use the largest size supported, up to 4MB. Under Solaris, a “super large pages” feature enables uses of pages up to 256MB. This feature is available for recent SPARC platforms. It can be enabled or disabled by using the `--super-large-pages` or `--skip-super-large-pages` option.

MySQL also supports the Linux implementation of large page support (which is called HugeTLB in Linux).

Before large pages can be used on Linux, the kernel must be enabled to support them and it is necessary to configure the HugeTLB memory pool. For reference, the HugeTBL API is documented in the `Documentation/vm/hugetlbpage.txt` file of your Linux sources.

The kernel for some recent systems such as Red Hat Enterprise Linux appear to have the large pages feature enabled by default. To check whether this is true for your kernel, use the following command and look for output lines containing “huge”:

```
shell> cat /proc/meminfo | grep -i huge
HugePages_Total:       0
HugePages_Free:        0
HugePages_Rsvd:        0
HugePages_Surp:        0
Hugepagesize:       4096 kB
```

The nonempty command output indicates that large page support is present, but the zero values indicate that no pages are configured for use.

If your kernel needs to be reconfigured to support large pages, consult the `hugetlbpage.txt` file for instructions.

Assuming that your Linux kernel has large page support enabled, configure it for use by MySQL using the following commands. Normally, you put these in an `rc` file or equivalent startup file that is executed during the system boot sequence, so that the commands execute each time the system starts. The commands should execute early in the boot sequence, before the MySQL server starts. Be sure to change the allocation numbers and the group number as appropriate for your system.

```
# Set the number of pages to be used.
# Each page is normally 2MB, so a value of 20 = 40MB.
# This command actually allocates memory, so this much
# memory must be available.
echo 20 > /proc/sys/vm/nr_hugepages

# Set the group number that is permitted to access this
# memory (102 in this case). The mysql user must be a
# member of this group.
echo 102 > /proc/sys/vm/hugetlb_shm_group

# Increase the amount of shmem permitted per segment
# (12G in this case).
echo 1560281088 > /proc/sys/kernel/shmmax

# Increase total amount of shared memory.  The value
# is the number of pages. At 4KB/page, 4194304 = 16GB.
echo 4194304 > /proc/sys/kernel/shmall
```

For MySQL usage, you normally want the value of `shmmax` to be close to the value of `shmall`.

To verify the large page configuration, check `/proc/meminfo` again as described previously. Now you should see some nonzero values:

```
shell> cat /proc/meminfo | grep -i huge
HugePages_Total:      20
HugePages_Free:       20
HugePages_Rsvd:        0
HugePages_Surp:        0
Hugepagesize:       4096 kB
```

The final step to make use of the `hugetlb_shm_group` is to give the `mysql` user an “unlimited” value for the memlock limit. This can be done either by editing `/etc/security/limits.conf` or by adding the following command to your `mysqld_safe` script:

```
ulimit -l unlimited
```

Adding the `ulimit` command to `mysqld_safe` causes the `root` user to set the memlock limit to `unlimited` before switching to the `mysql` user. (This assumes that `mysqld_safe` is started by `root`.)

Large page support in MySQL is disabled by default. To enable it, start the server with the `--large-pages` option. For example, you can use the following lines in your server's `my.cnf` file:

```
[mysqld]
large-pages
```

With this option, `InnoDB` uses large pages automatically for its buffer pool and additional memory pool. If `InnoDB` cannot do this, it falls back to use of traditional memory and writes a warning to the error log: `Warning: Using conventional memory pool`

To verify that large pages are being used, check `/proc/meminfo` again:

```
shell> cat /proc/meminfo | grep -i huge
HugePages_Total:      20
HugePages_Free:       20
HugePages_Rsvd:        2
```

```
HugePages_Surp:        0
Hugepagesize:       4096 kB
```

8.12.4 Optimizing Network Use

8.12.4.1 How MySQL Uses Threads for Client Connections

Connection manager threads handle client connection requests on the network interfaces that the server listens to. On all platforms, one manager thread handles TCP/IP connection requests. On Unix, this manager thread also handles Unix socket file connection requests. On Windows, a manager thread handles shared-memory connection requests, and another handles named-pipe connection requests. The server does not create threads to handle interfaces that it does not listen to. For example, a Windows server that does not have support for named-pipe connections enabled does not create a thread to handle them.

Connection manager threads associate each client connection with a thread dedicated to it that handles authentication and request processing for that connection. Manager threads create a new thread when necessary but try to avoid doing so by consulting the thread cache first to see whether it contains a thread that can be used for the connection. When a connection ends, its thread is returned to the thread cache if the cache is not full.

In this connection thread model, there are as many threads as there are clients currently connected, which has some disadvantages when server workload must scale to handle large numbers of connections. For example, thread creation and disposal becomes expensive. Also, each thread requires server and kernel resources, such as stack space. To accommodate a large number of simultaneous connections, the stack size per thread must be kept small, leading to a situation where it is either too small or the server consumes large amounts of memory. Exhaustion of other resources can occur as well, and scheduling overhead can become significant.

To control and monitor how the server manages threads that handle client connections, several system and status variables are relevant. (See Section 5.1.5, “Server System Variables”, and Section 5.1.7, “Server Status Variables”.)

The thread cache has a size determined by the `thread_cache_size` system variable. The default value is 0 (no caching), which causes a thread to be set up for each new connection and disposed of when the connection terminates. Set `thread_cache_size` to *`N`* to enable *`N`* inactive connection threads to be cached. `thread_cache_size` can be set at server startup or changed while the server runs. A connection thread becomes inactive when the client connection with which it was associated terminates.

To monitor the number of threads in the cache and how many threads have been created because a thread could not be taken from the cache, monitor the `Threads_cached` and `Threads_created` status variables.

You can set `max_connections` at server startup or at runtime to control the maximum number of clients that can connect simultaneously.

When the thread stack is too small, this limits the complexity of the SQL statements which the server can handle, the recursion depth of stored procedures, and other memory-consuming actions. To set a stack size of *`N`* bytes for each thread, start the server with `--thread_stack=`*`N`*.

8.12.4.2 DNS Lookup Optimization and the Host Cache

The MySQL server maintains a host cache in memory that contains information about clients: IP address, host name, and error information. The server uses this cache for nonlocal TCP connections. It does not use the cache for TCP connections established using a loopback interface address (`127.0.0.1` or `::1`), or for connections established using a Unix socket file, named pipe, or shared memory.

For each new client connection, the server uses the client IP address to check whether the client host name is in the host cache. If not, the server attempts to resolve the host name. First, it resolves the

IP address to a host name and resolves that host name back to an IP address. Then it compares the result to the original IP address to ensure that they are the same. The server stores information about the result of this operation in the host cache. If the cache is full, the least recently used entry is discarded.

The `host_cache` Performance Schema table exposes the contents of the host cache so that it can be examined using `SELECT` statements. This may help you diagnose the causes of connection problems. See Section 25.11.16.1, “The host_cache Table”.

The server handles entries in the host cache like this:

1. When the first TCP client connection reaches the server from a given IP address, a new entry is created to record the client IP, host name, and client lookup validation flag. Initially, the host name is set to `NULL` and the flag is false. This entry is also used for subsequent client connections from the same originating IP.

2. If the validation flag for the client IP entry is false, the server attempts an IP-to-host name DNS resolution. If that is successful, the host name is updated with the resolved host name and the validation flag is set to true. If resolution is unsuccessful, the action taken depends on whether the error is permanent or transient. For permanent failures, the host name remains `NULL` and the validation flag is set to true. For transient failures, the host name and validation flag remain unchanged. (Another DNS resolution attempt occurs the next time a client connects from this IP.)

3. If an error occurs while processing an incoming client connection from a given IP address, the server updates the corresponding error counters in the entry for that IP. For a description of the errors recorded, see Section 25.11.16.1, “The host_cache Table”.

The server performs host name resolution using the thread-safe `gethostbyaddr_r()` and `gethostbyname_r()` calls if the operating system supports them. Otherwise, the thread performing the lookup locks a mutex and calls `gethostbyaddr()` and `gethostbyname()` instead. In this case, no other thread can resolve host names that are not in the host cache until the thread holding the mutex lock releases it.

The server uses the host cache for several purposes:

- By caching the results of IP-to-host name lookups, the server avoids doing a DNS lookup for each client connection. Instead, for a given host, it needs to perform a lookup only for the first connection from that host.

- The cache contains information about errors that occur during the connection process. Some errors are considered “blocking.” If too many of these occur successively from a given host without a successful connection, the server blocks further connections from that host. The `max_connect_errors` system variable determines the number of permitted errors before blocking occurs. See Section B.5.2.5, “Host 'host_name' is blocked”.

To unblock blocked hosts, flush the host cache by issuing a `FLUSH HOSTS` statement or executing a `mysqladmin flush-hosts` command.

It is possible for a blocked host to become unblocked even without `FLUSH HOSTS` if activity from other hosts has occurred since the last connection attempt from the blocked host. This can occur because the server discards the least recently used cache entry to make room for a new entry if the cache is full when a connection arrives from a client IP not in the cache. If the discarded entry is for a blocked host, that host becomes unblocked.

The host cache is enabled by default. To disable it, set the `host_cache_size` system variable to 0, either at server startup or at runtime.

To disable DNS host name lookups, start the server with the `--skip-name-resolve` option. In this case, the server uses only IP addresses and not host names to match connecting hosts to rows in the MySQL grant tables. Only accounts specified in those tables using IP addresses can be used. (Be sure that an account exists that specifies an IP address or you may not be able to connect.)

If you have a very slow DNS and many hosts, you might be able to improve performance either by disabling DNS lookups with `--skip-name-resolve` or by increasing the value of `host_cache_size` to make the host cache larger.

To disallow TCP/IP connections entirely, start the server with the `--skip-networking` option.

Some connection errors are not associated with TCP connections, occur very early in the connection process (even before an IP address is known), or are not specific to any particular IP address (such as out-of-memory conditions). For information about these errors, check the `Connection_errors_xxx` status variables (see Section 5.1.7, “Server Status Variables”).

8.13 Measuring Performance (Benchmarking)

To measure performance, consider the following factors:

- Whether you are measuring the speed of a single operation on a quiet system, or how a set of operations (a “workload”) works over a period of time. With simple tests, you usually test how changing one aspect (a configuration setting, the set of indexes on a table, the SQL clauses in a query) affects performance. Benchmarks are typically long-running and elaborate performance tests, where the results could dictate high-level choices such as hardware and storage configuration, or how soon to upgrade to a new MySQL version.
- For benchmarking, sometimes you must simulate a heavy database workload to get an accurate picture.
- Performance can vary depending on so many different factors that a difference of a few percentage points might not be a decisive victory. The results might shift the opposite way when you test in a different environment.
- Certain MySQL features help or do not help performance depending on the workload. For completeness, always test performance with those features turned on and turned off. The two most important features to try with each workload are the MySQL query cache, and the adaptive hash index for `InnoDB` tables.

This section progresses from simple and direct measurement techniques that a single developer can do, to more complicated ones that require additional expertise to perform and interpret the results.

8.13.1 Measuring the Speed of Expressions and Functions

To measure the speed of a specific MySQL expression or function, invoke the `BENCHMARK()` function using the `mysql` client program. Its syntax is `BENCHMARK(loop_count,expression)`. The return value is always zero, but `mysql` prints a line displaying approximately how long the statement took to execute. For example:

```
mysql> SELECT BENCHMARK(1000000,1+1);
+------------------------+
| BENCHMARK(1000000,1+1) |
+------------------------+
|                      0 |
+------------------------+
1 row in set (0.32 sec)
```

This result was obtained on a Pentium II 400MHz system. It shows that MySQL can execute 1,000,000 simple addition expressions in 0.32 seconds on that system.

The built-in MySQL functions are typically highly optimized, but there may be some exceptions. `BENCHMARK()` is an excellent tool for finding out if some function is a problem for your queries.

8.13.2 Using Your Own Benchmarks

Benchmark your application and database to find out where the bottlenecks are. After fixing one bottleneck (or by replacing it with a “dummy” module), you can proceed to identify the next bottleneck. Even if the overall performance for your application currently is acceptable, you should at least make a plan for each bottleneck and decide how to solve it if someday you really need the extra performance.

A free benchmark suite is the Open Source Database Benchmark, available at http://osdb.sourceforge.net/.

It is very common for a problem to occur only when the system is very heavily loaded. We have had many customers who contact us when they have a (tested) system in production and have encountered load problems. In most cases, performance problems turn out to be due to issues of basic database design (for example, table scans are not good under high load) or problems with the operating system or libraries. Most of the time, these problems would be much easier to fix if the systems were not already in production.

To avoid problems like this, benchmark your whole application under the worst possible load:

- The `mysqlslap` program can be helpful for simulating a high load produced by multiple clients issuing queries simultaneously. See Section 4.5.8, “`mysqlslap` — Load Emulation Client”.
- You can also try benchmarking packages such as SysBench and DBT2, available at https://launchpad.net/sysbench, and http://osdldbt.sourceforge.net/#dbt2.

These programs or packages can bring a system to its knees, so be sure to use them only on your development systems.

8.13.3 Measuring Performance with performance_schema

You can query the tables in the `performance_schema` database to see real-time information about the performance characteristics of your server and the applications it is running. See Chapter 25, *MySQL Performance Schema* for details.

8.14 Examining Thread Information

When you are attempting to ascertain what your MySQL server is doing, it can be helpful to examine the process list, which is the set of threads currently executing within the server. Process list information is available from these sources:

- The `SHOW [FULL] PROCESSLIST` statement: Section 13.7.5.29, “SHOW PROCESSLIST Syntax”
- The `SHOW PROFILE` statement: Section 13.7.5.31, “SHOW PROFILES Syntax”
- The `INFORMATION_SCHEMA PROCESSLIST` table: Section 24.15, “The INFORMATION_SCHEMA PROCESSLIST Table”
- The `mysqladmin processlist` command: Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”
- The Performance Schema `threads` table, stage tables, and lock tables: Section 25.11.16, “Performance Schema Miscellaneous Tables”, Section 25.11.5, “Performance Schema Stage Event Tables”, Section 25.11.12, “Performance Schema Lock Tables”.

Access to `threads` does not require a mutex and has minimal impact on server performance. `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` have negative performance consequences because they require a mutex. `threads` also shows information about background threads, which `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` do not. This means that `threads` can be used to monitor activity the other thread information sources cannot.

You can always view information about your own threads. To view information about threads being executed for other accounts, you must have the `PROCESS` privilege.

Each process list entry contains several pieces of information:

- `Id` is the connection identifier for the client associated with the thread.
- `User` and `Host` indicate the account associated with the thread.
- `db` is the default database for the thread, or `NULL` if none is selected.
- `Command` and `State` indicate what the thread is doing.

  Most states correspond to very quick operations. If a thread stays in a given state for many seconds, there might be a problem that needs to be investigated.
- `Time` indicates how long the thread has been in its current state. The thread's notion of the current time may be altered in some cases: The thread can change the time with `SET TIMESTAMP = value`. For a thread running on a slave that is processing events from the master, the thread time is set to the time found in the events and thus reflects current time on the master and not the slave.
- `Info` contains the text of the statement being executed by the thread, or `NULL` if it is not executing one. By default, this value contains only the first 100 characters of the statement. To see the complete statements, use `SHOW FULL PROCESSLIST`.
- The `sys` schema `processlist` view, which presents information from the Performance Schema `threads` table in a more accessible format: Section 26.4.3.22, “The processlist and x$processlist Views”
- The `sys` schema `session` view, which presents information about user sessions (like the `sys` schema `processlist` view, but with background processes filtered out): Section 26.4.3.33, “The session and x$session Views”

The following sections list the possible `Command` values, and `State` values grouped by category. The meaning for some of these values is self-evident. For others, additional description is provided.

8.14.1 Thread Command Values

A thread can have any of the following `Command` values:

- `Binlog Dump`

  This is a thread on a master server for sending binary log contents to a slave server.
- `Change user`

  The thread is executing a change-user operation.
- `Close stmt`

  The thread is closing a prepared statement.
- `Connect`

  A replication slave is connected to its master.
- `Connect Out`

  A replication slave is connecting to its master.
- `Create DB`

  The thread is executing a create-database operation.
- `Daemon`

  This thread is internal to the server, not a thread that services a client connection.
- `Debug`

The thread is generating debugging information.

- `Delayed insert`

  The thread is a delayed-insert handler.

- `Drop DB`

  The thread is executing a drop-database operation.

- `Error`

- `Execute`

  The thread is executing a prepared statement.

- `Fetch`

  The thread is fetching the results from executing a prepared statement.

- `Field List`

  The thread is retrieving information for table columns.

- `Init DB`

  The thread is selecting a default database.

- `Kill`

  The thread is killing another thread.

- `Long Data`

  The thread is retrieving long data in the result of executing a prepared statement.

- `Ping`

  The thread is handling a server-ping request.

- `Prepare`

  The thread is preparing a prepared statement.

- `Processlist`

  The thread is producing information about server threads.

- `Query`

  The thread is executing a statement.

- `Quit`

  The thread is terminating.

- `Refresh`

  The thread is flushing table, logs, or caches, or resetting status variable or replication server information.

- `Register Slave`

  The thread is registering a slave server.

- `Reset stmt`

  The thread is resetting a prepared statement.

- `Set option`

  The thread is setting or resetting a client statement-execution option.

- `Shutdown`

  The thread is shutting down the server.

- `Sleep`

  The thread is waiting for the client to send a new statement to it.

- `Statistics`

  The thread is producing server-status information.

- `Table Dump`

  The thread is sending table contents to a slave server.

- `Time`

  Unused.

8.14.2 General Thread States

The following list describes thread `State` values that are associated with general query processing and not more specialized activities such as replication. Many of these are useful only for finding bugs in the server.

- `After create`

  This occurs when the thread creates a table (including internal temporary tables), at the end of the function that creates the table. This state is used even if the table could not be created due to some error.

- `Analyzing`

  The thread is calculating a `MyISAM` table key distributions (for example, for `ANALYZE TABLE`).

- `checking permissions`

  The thread is checking whether the server has the required privileges to execute the statement.

- `Checking table`

  The thread is performing a table check operation.

- `cleaning up`

  The thread has processed one command and is preparing to free memory and reset certain state variables.

- `closing tables`

  The thread is flushing the changed table data to disk and closing the used tables. This should be a fast operation. If not, verify that you do not have a full disk and that the disk is not in very heavy use.

- `converting HEAP to MyISAM`

The thread is converting an internal temporary table from a `MEMORY` table to an on-disk `MyISAM` table.

- `copy to tmp table`

  The thread is processing an `ALTER TABLE` statement. This state occurs after the table with the new structure has been created but before rows are copied into it.

  For a thread in this state, the Performance Schema can be used to obtain about the progress of the copy operation. See Section 25.11.5, “Performance Schema Stage Event Tables”.

- `Copying to group table`

  If a statement has different `ORDER BY` and `GROUP BY` criteria, the rows are sorted by group and copied to a temporary table.

- `Copying to tmp table`

  The server is copying to a temporary table in memory.

- `altering table`

  The server is in the process of executing an in-place `ALTER TABLE`.

- `Copying to tmp table on disk`

  The server is copying to a temporary table on disk. The temporary result set has become too large (see Section 8.4.4, “Internal Temporary Table Use in MySQL”). Consequently, the thread is changing the temporary table from in-memory to disk-based format to save memory.

- `Creating index`

  The thread is processing `ALTER TABLE ... ENABLE KEYS` for a `MyISAM` table.

- `Creating sort index`

  The thread is processing a `SELECT` that is resolved using an internal temporary table.

- `creating table`

  The thread is creating a table. This includes creation of temporary tables.

- `Creating tmp table`

  The thread is creating a temporary table in memory or on disk. If the table is created in memory but later is converted to an on-disk table, the state during that operation will be `Copying to tmp table on disk`.

- `committing alter table to storage engine`

  The server has finished an in-place `ALTER TABLE` and is committing the result.

- `deleting from main table`

  The server is executing the first part of a multiple-table delete. It is deleting only from the first table, and saving columns and offsets to be used for deleting from the other (reference) tables.

- `deleting from reference tables`

  The server is executing the second part of a multiple-table delete and deleting the matched rows from the other tables.

- `discard_or_import_tablespace`

The thread is processing an `ALTER TABLE ... DISCARD TABLESPACE` or `ALTER TABLE ... IMPORT TABLESPACE` statement.

- `end`

  This occurs at the end but before the cleanup of `ALTER TABLE`, `CREATE VIEW`, `DELETE`, `INSERT`, `SELECT`, or `UPDATE` statements.

- `executing`

  The thread has begun executing a statement.

- `Execution of init_command`

  The thread is executing statements in the value of the `init_command` system variable.

- `freeing items`

  The thread has executed a command. Some freeing of items done during this state involves the query cache. This state is usually followed by `cleaning up`.

- `FULLTEXT initialization`

  The server is preparing to perform a natural-language full-text search.

- `init`

  This occurs before the initialization of `ALTER TABLE`, `DELETE`, `INSERT`, `SELECT`, or `UPDATE` statements. Actions taken by the server in this state include flushing the binary log, the `InnoDB` log, and some query cache cleanup operations.

  For the `end` state, the following operations could be happening:

  - Removing query cache entries after data in a table is changed
  - Writing an event to the binary log
  - Freeing memory buffers, including for blobs

- `Killed`

  Someone has sent a `KILL` statement to the thread and it should abort next time it checks the kill flag. The flag is checked in each major loop in MySQL, but in some cases it might still take a short time for the thread to die. If the thread is locked by some other thread, the kill takes effect as soon as the other thread releases its lock.

- `Locking system tables`

  The thread is trying to lock a system table (for example, a time zone or log table).

- `logging slow query`

  The thread is writing a statement to the slow-query log.

- `login`

  The initial state for a connection thread until the client has been authenticated successfully.

- `manage keys`

  The server is enabling or disabling a table index.

- `NULL`

This state is used for the `SHOW PROCESSLIST` state.

- `Opening system tables`

  The thread is trying to open a system table (for example, a time zone or log table).

- `Opening tables`

  The thread is trying to open a table. This is should be very fast procedure, unless something prevents opening. For example, an `ALTER TABLE` or a `LOCK TABLE` statement can prevent opening a table until the statement is finished. It is also worth checking that your `table_open_cache` value is large enough.

  For system tables, the `Opening system tables` state is used instead.

- `optimizing`

  The server is performing initial optimizations for a query.

- `preparing`

  This state occurs during query optimization.

- `Purging old relay logs`

  The thread is removing unneeded relay log files.

- `query end`

  This state occurs after processing a query but before the `freeing items` state.

- `Receiving from client`

  The server is reading a packet from the client.

- `Removing duplicates`

  The query was using `SELECT DISTINCT` in such a way that MySQL could not optimize away the distinct operation at an early stage. Because of this, MySQL requires an extra stage to remove all duplicated rows before sending the result to the client.

- `removing tmp table`

  The thread is removing an internal temporary table after processing a `SELECT` statement. This state is not used if no temporary table was created.

- `rename`

  The thread is renaming a table.

- `rename result table`

  The thread is processing an `ALTER TABLE` statement, has created the new table, and is renaming it to replace the original table.

- `Reopen tables`

  The thread got a lock for the table, but noticed after getting the lock that the underlying table structure changed. It has freed the lock, closed the table, and is trying to reopen it.

- `Repair by sorting`

  The repair code is using a sort to create indexes.

- `preparing for alter table`

  The server is preparing to execute an in-place `ALTER TABLE`.

- `Repair done`

  The thread has completed a multi-threaded repair for a `MyISAM` table.

- `Repair with keycache`

  The repair code is using creating keys one by one through the key cache. This is much slower than `Repair by sorting`.

- `Rolling back`

  The thread is rolling back a transaction.

- `Saving state`

  For `MyISAM` table operations such as repair or analysis, the thread is saving the new table state to the `.MYI` file header. State includes information such as number of rows, the `AUTO_INCREMENT` counter, and key distributions.

- `Searching rows for update`

  The thread is doing a first phase to find all matching rows before updating them. This has to be done if the `UPDATE` is changing the index that is used to find the involved rows.

- `Sending data`

  The thread is reading and processing rows for a `SELECT` statement, and sending data to the client. Because operations occurring during this state tend to perform large amounts of disk access (reads), it is often the longest-running state over the lifetime of a given query.

- `Sending to client`

  The server is writing a packet to the client.

- `setup`

  The thread is beginning an `ALTER TABLE` operation.

- `Sorting for group`

  The thread is doing a sort to satisfy a `GROUP BY`.

- `Sorting for order`

  The thread is doing a sort to satisfy an `ORDER BY`.

- `Sorting index`

  The thread is sorting index pages for more efficient access during a `MyISAM` table optimization operation.

- `Sorting result`

  For a `SELECT` statement, this is similar to `Creating sort index`, but for nontemporary tables.

- `statistics`

  The server is calculating statistics to develop a query execution plan. If a thread is in this state for a long time, the server is probably disk-bound performing other work.

- `System lock`

  The thread has called `mysql_lock_tables()` and the thread state has not been updated since. This is a very general state that can occur for many reasons.

  For example, the thread is going to request or is waiting for an internal or external system lock for the table. This can occur when `InnoDB` waits for a table-level lock during execution of `LOCK TABLES`. If this state is being caused by requests for external locks and you are not using multiple `mysqld` servers that are accessing the same `MyISAM` tables, you can disable external system locks with the `--skip-external-locking` option. However, external locking is disabled by default, so it is likely that this option will have no effect. For `SHOW PROFILE`, this state means the thread is requesting the lock (not waiting for it).

  For system tables, the `Locking system tables` state is used instead.

- `update`

  The thread is getting ready to start updating the table.

- `Updating`

  The thread is searching for rows to update and is updating them.

- `updating main table`

  The server is executing the first part of a multiple-table update. It is updating only the first table, and saving columns and offsets to be used for updating the other (reference) tables.

- `updating reference tables`

  The server is executing the second part of a multiple-table update and updating the matched rows from the other tables.

- `User lock`

  The thread is going to request or is waiting for an advisory lock requested with a `GET_LOCK()` call. For `SHOW PROFILE`, this state means the thread is requesting the lock (not waiting for it).

- `User sleep`

  The thread has invoked a `SLEEP()` call.

- `Waiting for commit lock`

  `FLUSH TABLES WITH READ LOCK` is waiting for a commit lock.

- `Waiting for global read lock`

  `FLUSH TABLES WITH READ LOCK` is waiting for a global read lock or the global `read_only` system variable is being set.

- `Waiting for tables`

  The thread got a notification that the underlying structure for a table has changed and it needs to reopen the table to get the new structure. However, to reopen the table, it must wait until all other threads have closed the table in question.

  This notification takes place if another thread has used `FLUSH TABLES` or one of the following statements on the table in question: `FLUSH TABLES` *`tbl_name`*, `ALTER TABLE`, `RENAME TABLE`, `REPAIR TABLE`, `ANALYZE TABLE`, or `OPTIMIZE TABLE`.

- `Waiting for table flush`

The thread is executing `FLUSH TABLES` and is waiting for all threads to close their tables, or the thread got a notification that the underlying structure for a table has changed and it needs to reopen the table to get the new structure. However, to reopen the table, it must wait until all other threads have closed the table in question.

This notification takes place if another thread has used `FLUSH TABLES` or one of the following statements on the table in question: `FLUSH TABLES tbl_name`, `ALTER TABLE`, `RENAME TABLE`, `REPAIR TABLE`, `ANALYZE TABLE`, or `OPTIMIZE TABLE`.

- `Waiting for lock_type lock`

  The server is waiting to acquire a `THR_LOCK` lock or a lock from the metadata locking subsystem, where `lock_type` indicates the type of lock.

  This state indicates a wait for a `THR_LOCK`:

  - `Waiting for table level lock`

  These states indicate a wait for a metadata lock:

  - `Waiting for event metadata lock`
  - `Waiting for global read lock`
  - `Waiting for schema metadata lock`
  - `Waiting for stored function metadata lock`
  - `Waiting for stored procedure metadata lock`
  - `Waiting for table metadata lock`
  - `Waiting for trigger metadata lock`

  For information about table lock indicators, see Section 8.11.1, “Internal Locking Methods”. For information about metadata locking, see Section 8.11.4, “Metadata Locking”. To see which locks are blocking lock requests, use the Performance Schema lock tables described at Section 25.11.12, “Performance Schema Lock Tables”.

- `Waiting on cond`

  A generic state in which the thread is waiting for a condition to become true. No specific state information is available.

- `Writing to net`

  The server is writing a packet to the network.

8.14.3 Query Cache Thread States

These thread states are associated with the query cache (see Section 8.10.3, “The MySQL Query Cache”).

- `checking privileges on cached query`

  The server is checking whether the user has privileges to access a cached query result.

- `checking query cache for query`

  The server is checking whether the current query is present in the query cache.

- `invalidating query cache entries`

Query cache entries are being marked invalid because the underlying tables have changed.

- `sending cached result to client`

  The server is taking the result of a query from the query cache and sending it to the client.

- `storing result in query cache`

  The server is storing the result of a query in the query cache.

- `Waiting for query cache lock`

  This state occurs while a session is waiting to take the query cache lock. This can happen for any statement that needs to perform some query cache operation, such as an `INSERT` or `DELETE` that invalidates the query cache, a `SELECT` that looks for a cached entry, `RESET QUERY CACHE`, and so forth.

8.14.4 Replication Master Thread States

The following list shows the most common states you may see in the `State` column for the master's `Binlog Dump` thread. If you see no `Binlog Dump` threads on a master server, this means that replication is not running; that is, that no slaves are currently connected.

- `Finished reading one binlog; switching to next binlog`

  The thread has finished reading a binary log file and is opening the next one to send to the slave.

- `Master has sent all binlog to slave; waiting for more updates`

  The thread has read all remaining updates from the binary logs and sent them to the slave. The thread is now idle, waiting for new events to appear in the binary log resulting from new updates occurring on the master.

- `Sending binlog event to slave`

  Binary logs consist of *events*, where an event is usually an update plus some other information. The thread has read an event from the binary log and is now sending it to the slave.

- `Waiting to finalize termination`

  A very brief state that occurs as the thread is stopping.

8.14.5 Replication Slave I/O Thread States

The following list shows the most common states you see in the `State` column for a slave server I/O thread. This state also appears in the `Slave_IO_State` column displayed by `SHOW SLAVE STATUS`, so you can get a good view of what is happening by using that statement.

- `Checking master version`

  A state that occurs very briefly, after the connection to the master is established.

- `Connecting to master`

  The thread is attempting to connect to the master.

- `Queueing master event to the relay log`

  The thread has read an event and is copying it to the relay log so that the SQL thread can process it.

- `Reconnecting after a failed binlog dump request`

The thread is trying to reconnect to the master.

- `Reconnecting after a failed master event read`

  The thread is trying to reconnect to the master. When connection is established again, the state becomes `Waiting for master to send event`.

- `Registering slave on master`

  A state that occurs very briefly after the connection to the master is established.

- `Requesting binlog dump`

  A state that occurs very briefly, after the connection to the master is established. The thread sends to the master a request for the contents of its binary logs, starting from the requested binary log file name and position.

- `Waiting for its turn to commit`

  A state that occurs when the slave thread is waiting for older worker threads to commit if `slave_preserve_commit_order` is enabled.

- `Waiting for master to send event`

  The thread has connected to the master and is waiting for binary log events to arrive. This can last for a long time if the master is idle. If the wait lasts for `slave_net_timeout` seconds, a timeout occurs. At that point, the thread considers the connection to be broken and makes an attempt to reconnect.

- `Waiting for master update`

  The initial state before `Connecting to master`.

- `Waiting for slave mutex on exit`

  A state that occurs briefly as the thread is stopping.

- `Waiting for the slave SQL thread to free enough relay log space`

  You are using a nonzero `relay_log_space_limit` value, and the relay logs have grown large enough that their combined size exceeds this value. The I/O thread is waiting until the SQL thread frees enough space by processing relay log contents so that it can delete some relay log files.

- `Waiting to reconnect after a failed binlog dump request`

  If the binary log dump request failed (due to disconnection), the thread goes into this state while it sleeps, then tries to reconnect periodically. The interval between retries can be specified using the `CHANGE MASTER TO` statement.

- `Waiting to reconnect after a failed master event read`

  An error occurred while reading (due to disconnection). The thread is sleeping for the number of seconds set by the `CHANGE MASTER TO` statement (default 60) before attempting to reconnect.

8.14.6 Replication Slave SQL Thread States

The following list shows the most common states you may see in the `State` column for a slave server SQL thread:

- `Killing slave`

  The thread is processing a `STOP SLAVE` statement.

- `Making temporary file (append) before replaying LOAD DATA INFILE`

  The thread is executing a `LOAD DATA INFILE` statement and is appending the data to a temporary file containing the data from which the slave will read rows.

- `Making temporary file (create) before replaying LOAD DATA INFILE`

  The thread is executing a `LOAD DATA INFILE` statement and is creating a temporary file containing the data from which the slave will read rows. This state can only be encountered if the original `LOAD DATA INFILE` statement was logged by a master running a version of MySQL earlier than version 5.0.3.

- `Reading event from the relay log`

  The thread has read an event from the relay log so that the event can be processed.

- `Slave has read all relay log; waiting for more updates`

  The thread has processed all events in the relay log files, and is now waiting for the I/O thread to write new events to the relay log.

- `Waiting for an event from Coordinator`

  Using the multi-threaded slave (`slave_parallel_workers` is greater than 1), one of the slave worker threads is waiting for an event from the coordinator thread.

- `Waiting for slave mutex on exit`

  A very brief state that occurs as the thread is stopping.

- `Waiting for Slave Workers to free pending events`

  This waiting action occurs when the total size of events being processed by Workers exceeds the size of the `slave_pending_jobs_size_max` system variable. The Coordinator resumes scheduling when the size drops below this limit. This state occurs only when `slave_parallel_workers` is set greater than 0.

- `Waiting for the next event in relay log`

  The initial state before `Reading event from the relay log`.

- `Waiting until MASTER_DELAY seconds after master executed event`

  The SQL thread has read an event but is waiting for the slave delay to lapse. This delay is set with the `MASTER_DELAY` option of `CHANGE MASTER TO`.

The `Info` column for the SQL thread may also show the text of a statement. This indicates that the thread has read an event from the relay log, extracted the statement from it, and may be executing it.

8.14.7 Replication Slave Connection Thread States

These thread states occur on a replication slave but are associated with connection threads, not with the I/O or SQL threads.

- `Changing master`

  The thread is processing a `CHANGE MASTER TO` statement.

- `Killing slave`

  The thread is processing a `STOP SLAVE` statement.

- `Opening master dump table`

This state occurs after `Creating table from master dump`.

- `Reading master dump table data`

  This state occurs after `Opening master dump table`.

- `Rebuilding the index on master dump table`

  This state occurs after `Reading master dump table data`.

8.14.8 Event Scheduler Thread States

These states occur for the Event Scheduler thread, threads that are created to execute scheduled events, or threads that terminate the scheduler.

- `Clearing`

  The scheduler thread or a thread that was executing an event is terminating and is about to end.

- `Initialized`

  The scheduler thread or a thread that will execute an event has been initialized.

- `Waiting for next activation`

  The scheduler has a nonempty event queue but the next activation is in the future.

- `Waiting for scheduler to stop`

  The thread issued `SET GLOBAL event_scheduler=OFF` and is waiting for the scheduler to stop.

- `Waiting on empty queue`

  The scheduler's event queue is empty and it is sleeping.

Chapter 9 Language Structure

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This chapter discusses the rules for writing the following elements of SQL statements when using MySQL:

- Literal values such as strings and numbers
- Identifiers such as database, table, and column names
- Keywords and reserved words
- User-defined and system variables
- Comments

9.1 Literal Values

This section describes how to write literal values in MySQL. These include strings, numbers, hexadecimal and bit values, boolean values, and `NULL`. The section also covers various nuances that you may encounter when dealing with these basic types in MySQL.

9.1.1 String Literals

A string is a sequence of bytes or characters, enclosed within either single quote (`'`) or double quote (`"`) characters. Examples:

```
'a string'
"another string"
```

Quoted strings placed next to each other are concatenated to a single string. The following lines are equivalent:

```
'a string'
'a' ' ' 'string'
```

If the `ANSI_QUOTES` SQL mode is enabled, string literals can be quoted only within single quotation marks because a string quoted within double quotation marks is interpreted as an identifier.

A *binary string* is a string of bytes. Every binary string has a character set and collation named `binary`. A *nonbinary string* is a string of characters. It has a character set other than `binary` and a collation that is compatible with the character set.

For both types of strings, comparisons are based on the numeric values of the string unit. For binary strings, the unit is the byte; comparisons use numeric byte values. For nonbinary strings, the unit is the character and some character sets support multibyte characters; comparisons use numeric character code values. Character code ordering is a function of the string collation. (For more information, see Section 10.1.8.5, “The binary Collation Compared to _bin Collations”.)

A character string literal may have an optional character set introducer and `COLLATE` clause, to designate it as a string that uses a particular character set and collation:

```
[_charset_name]'string' [COLLATE collation_name]
```

Examples:

```
SELECT _latin1'string';
SELECT _binary'string';
SELECT _utf8'string' COLLATE utf8_danish_ci;
```

You can use `N'literal'` (or `n'literal'`) to create a string in the national character set. These statements are equivalent:

```
SELECT N'some text';
SELECT n'some text';
SELECT _utf8'some text';
```

For information about these forms of string syntax, see Section 10.1.3.7, “The National Character Set”, and Section 10.1.3.8, “Character Set Introducers”.

Within a string, certain sequences have special meaning unless the `NO_BACKSLASH_ESCAPES` SQL mode is enabled. Each of these sequences begins with a backslash (`\`), known as the *escape character*. MySQL recognizes the escape sequences shown in Table 9.1, “Special Character Escape Sequences”. For all other escape sequences, backslash is ignored. That is, the escaped character is interpreted as if it was not escaped. For example, `\x` is just `x`. These sequences are case sensitive. For example, `\b` is interpreted as a backspace, but `\B` is interpreted as `B`. Escape processing is done according to the character set indicated by the `character_set_connection` system variable. This is true even for strings that are preceded by an introducer that indicates a different character set, as discussed in Section 10.1.3.6, “Character String Literal Character Set and Collation”.

Table 9.1 Special Character Escape Sequences

Escape Sequence	Character Represented by Sequence
`\0`	An ASCII NUL (`X'00'`) character
`\'`	A single quote (`'`) character
`\"`	A double quote (`"`) character
`\b`	A backspace character
`\n`	A newline (linefeed) character
`\r`	A carriage return character
`\t`	A tab character
`\Z`	ASCII 26 (Control+Z); see note following the table
`\\`	A backslash (`\`) character
`\%`	A `%` character; see note following the table
`\_`	A `_` character; see note following the table

The ASCII 26 character can be encoded as `\Z` to enable you to work around the problem that ASCII 26 stands for END-OF-FILE on Windows. ASCII 26 within a file causes problems if you try to use `mysql db_name < file_name`.

The `\%` and `\_` sequences are used to search for literal instances of `%` and `_` in pattern-matching contexts where they would otherwise be interpreted as wildcard characters. See the description of the `LIKE` operator in Section 12.5.1, “String Comparison Functions”. If you use `\%` or `\_` outside of pattern-matching contexts, they evaluate to the strings `\%` and `\_`, not to `%` and `_`.

There are several ways to include quote characters within a string:

- A `'` inside a string quoted with `'` may be written as `''`.
- A `"` inside a string quoted with `"` may be written as `""`.
- Precede the quote character by an escape character (`\`).
- A `'` inside a string quoted with `"` needs no special treatment and need not be doubled or escaped. In the same way, `"` inside a string quoted with `'` needs no special treatment.

The following `SELECT` statements demonstrate how quoting and escaping work:

```
mysql> SELECT 'hello', '"hello"', '""hello""', 'hel''lo', '\'hello';
+-------+---------+-----------+--------+--------+
| hello | "hello" | ""hello"" | hel'lo | 'hello |
+-------+---------+-----------+--------+--------+

mysql> SELECT "hello", "'hello'", "''hello''", "hel""lo", "\"hello";
+-------+---------+-----------+--------+--------+
| hello | 'hello' | ''hello'' | hel"lo | "hello |
+-------+---------+-----------+--------+--------+

mysql> SELECT 'This\nIs\nFour\nLines';
+--------------------+
| This
Is
Four
Lines |
+--------------------+

mysql> SELECT 'disappearing\ backslash';
+------------------------+
| disappearing backslash |
+------------------------+
```

To insert binary data into a string column (such as a `BLOB` column), you should represent certain characters by escape sequences. Backslash (`\`) and the quote character used to quote the string must be escaped. In certain client environments, it may also be necessary to escape `NUL` or Control+Z. The `mysql` client truncates quoted strings containing `NUL` characters if they are not escaped, and Control+Z may be taken for END-OF-FILE on Windows if not escaped. For the escape sequences that represent each of these characters, see Table 9.1, “Special Character Escape Sequences”.

When writing application programs, any string that might contain any of these special characters must be properly escaped before the string is used as a data value in an SQL statement that is sent to the MySQL server. You can do this in two ways:

- Process the string with a function that escapes the special characters. In a C program, you can use the `mysql_real_escape_string_quote()` C API function to escape characters. See Section 27.7.7.56, “mysql_real_escape_string_quote()”. Within SQL statements that construct other SQL statements, you can use the `QUOTE()` function. The Perl DBI interface provides a `quote` method to convert special characters to the proper escape sequences. See Section 27.9, “MySQL Perl API”. Other language interfaces may provide a similar capability.
- As an alternative to explicitly escaping special characters, many MySQL APIs provide a placeholder capability that enables you to insert special markers into a statement string, and then bind data

values to them when you issue the statement. In this case, the API takes care of escaping special characters in the values for you.

9.1.2 Numeric Literals

Number literals include exact-value (integer and `DECIMAL`) literals and approximate-value (floating-point) literals.

Integers are represented as a sequence of digits. Numbers may include `.` as a decimal separator. Numbers may be preceded by `-` or `+` to indicate a negative or positive value, respectively. Numbers represented in scientific notation with a mantissa and exponent are approximate-value numbers.

Exact-value numeric literals have an integer part or fractional part, or both. They may be signed. Examples: `1`, `.2`, `3.4`, `-5`, `-6.78`, `+9.10`.

Approximate-value numeric literals are represented in scientific notation with a mantissa and exponent. Either or both parts may be signed. Examples: `1.2E3`, `1.2E-3`, `-1.2E3`, `-1.2E-3`.

Two numbers that look similar may be treated differently. For example, `2.34` is an exact-value (fixed-point) number, whereas `2.34E0` is an approximate-value (floating-point) number.

The `DECIMAL` data type is a fixed-point type and calculations are exact. In MySQL, the `DECIMAL` type has several synonyms: `NUMERIC`, `DEC`, `FIXED`. The integer types also are exact-value types. For more information about exact-value calculations, see Section 12.21, “Precision Math”.

The `FLOAT` and `DOUBLE` data types are floating-point types and calculations are approximate. In MySQL, types that are synonymous with `FLOAT` or `DOUBLE` are `DOUBLE PRECISION` and `REAL`.

An integer may be used in a floating-point context; it is interpreted as the equivalent floating-point number.

9.1.3 Date and Time Literals

Date and time values can be represented in several formats, such as quoted strings or as numbers, depending on the exact type of the value and other factors. For example, in contexts where MySQL expects a date, it interprets any of `'2015-07-21'`, `'20150721'`, and `20150721` as a date.

This section describes the acceptable formats for date and time literals. For more information about the temporal data types, such as the range of permitted values, consult these sections:

- Section 11.1.2, “Date and Time Type Overview”
- Section 11.3, “Date and Time Types”

Standard SQL and ODBC Date and Time Literals. Standard SQL permits temporal literals to be specified using a type keyword and a string. The space between the keyword and string is optional.

```
DATE 'str'
TIME 'str'
TIMESTAMP 'str'
```

MySQL recognizes those constructions and also the corresponding ODBC syntax:

```
{ d 'str' }
{ t 'str' }
{ ts 'str' }
```

MySQL uses the type keyword and these constructions produce `DATE`, `TIME`, and `DATETIME` values, respectively, including a trailing fractional seconds part if specified. The `TIMESTAMP` syntax produces a `DATETIME` value in MySQL because `DATETIME` has a range that more closely corresponds to the standard SQL `TIMESTAMP` type, which has a year range from `0001` to `9999`. (The MySQL `TIMESTAMP` year range is `1970` to `2038`.)

String and Numeric Literals in Date and Time Context. MySQL recognizes `DATE` values in these formats:

- As a string in either `'YYYY-MM-DD'` or `'YY-MM-DD'` format. A “relaxed” syntax is permitted: Any punctuation character may be used as the delimiter between date parts. For example, `'2012-12-31'`, `'2012/12/31'`, `'2012^12^31'`, and `'2012@12@31'` are equivalent.
- As a string with no delimiters in either `'YYYYMMDD'` or `'YYMMDD'` format, provided that the string makes sense as a date. For example, `'20070523'` and `'070523'` are interpreted as `'2007-05-23'`, but `'071332'` is illegal (it has nonsensical month and day parts) and becomes `'0000-00-00'`.
- As a number in either `YYYYMMDD` or `YYMMDD` format, provided that the number makes sense as a date. For example, `19830905` and `830905` are interpreted as `'1983-09-05'`.

MySQL recognizes `DATETIME` and `TIMESTAMP` values in these formats:

- As a string in either `'YYYY-MM-DD HH:MM:SS'` or `'YY-MM-DD HH:MM:SS'` format. A “relaxed” syntax is permitted here, too: Any punctuation character may be used as the delimiter between date parts or time parts. For example, `'2012-12-31 11:30:45'`, `'2012^12^31 11+30+45'`, `'2012/12/31 11*30*45'`, and `'2012@12@31 11^30^45'` are equivalent.

  The only delimiter recognized between a date and time part and a fractional seconds part is the decimal point.

  The date and time parts can be separated by `T` rather than a space. For example, `'2012-12-31 11:30:45'` `'2012-12-31T11:30:45'` are equivalent.
- As a string with no delimiters in either `'YYYYMMDDHHMMSS'` or `'YYMMDDHHMMSS'` format, provided that the string makes sense as a date. For example, `'20070523091528'` and `'070523091528'` are interpreted as `'2007-05-23 09:15:28'`, but `'071122129015'` is illegal (it has a nonsensical minute part) and becomes `'0000-00-00 00:00:00'`.
- As a number in either `YYYYMMDDHHMMSS` or `YYMMDDHHMMSS` format, provided that the number makes sense as a date. For example, `19830905132800` and `830905132800` are interpreted as `'1983-09-05 13:28:00'`.

A `DATETIME` or `TIMESTAMP` value can include a trailing fractional seconds part in up to microseconds (6 digits) precision. The fractional part should always be separated from the rest of the time by a decimal point; no other fractional seconds delimiter is recognized. For information about fractional seconds support in MySQL, see Section 11.3.6, “Fractional Seconds in Time Values”.

Dates containing two-digit year values are ambiguous because the century is unknown. MySQL interprets two-digit year values using these rules:

- Year values in the range `70-99` are converted to `1970-1999`.
- Year values in the range `00-69` are converted to `2000-2069`.

See also Section 11.3.8, “Two-Digit Years in Dates”.

For values specified as strings that include date part delimiters, it is unnecessary to specify two digits for month or day values that are less than `10`. `'2015-6-9'` is the same as `'2015-06-09'`. Similarly, for values specified as strings that include time part delimiters, it is unnecessary to specify two digits for hour, minute, or second values that are less than `10`. `'2015-10-30 1:2:3'` is the same as `'2015-10-30 01:02:03'`.

Values specified as numbers should be 6, 8, 12, or 14 digits long. If a number is 8 or 14 digits long, it is assumed to be in `YYYYMMDD` or `YYYYMMDDHHMMSS` format and that the year is given by the first 4 digits. If the number is 6 or 12 digits long, it is assumed to be in `YYMMDD` or `YYMMDDHHMMSS` format and that the year is given by the first 2 digits. Numbers that are not one of these lengths are interpreted as though padded with leading zeros to the closest length.

Values specified as nondelimited strings are interpreted according their length. For a string 8 or 14 characters long, the year is assumed to be given by the first 4 characters. Otherwise, the year is assumed to be given by the first 2 characters. The string is interpreted from left to right to find year, month, day, hour, minute, and second values, for as many parts as are present in the string. This means you should not use strings that have fewer than 6 characters. For example, if you specify `'9903'`, thinking that represents March, 1999, MySQL converts it to the "zero" date value. This occurs because the year and month values are `99` and `03`, but the day part is completely missing. However, you can explicitly specify a value of zero to represent missing month or day parts. For example, to insert the value `'1999-03-00'`, use `'990300'`.

MySQL recognizes `TIME` values in these formats:

- As a string in `'D HH:MM:SS'` format. You can also use one of the following "relaxed" syntaxes: `'HH:MM:SS'`, `'HH:MM'`, `'D HH:MM'`, `'D HH'`, or `'SS'`. Here `D` represents days and can have a value from 0 to 34.
- As a string with no delimiters in `'HHMMSS'` format, provided that it makes sense as a time. For example, `'101112'` is understood as `'10:11:12'`, but `'109712'` is illegal (it has a nonsensical minute part) and becomes `'00:00:00'`.
- As a number in `HHMMSS` format, provided that it makes sense as a time. For example, `101112` is understood as `'10:11:12'`. The following alternative formats are also understood: `SS`, `MMSS`, or `HHMMSS`.

A trailing fractional seconds part is recognized in the `'D HH:MM:SS.fraction'`, `'HH:MM:SS.fraction'`, `'HHMMSS.fraction'`, and `HHMMSS.fraction` time formats, where `fraction` is the fractional part in up to microseconds (6 digits) precision. The fractional part should always be separated from the rest of the time by a decimal point; no other fractional seconds delimiter is recognized. For information about fractional seconds support in MySQL, see Section 11.3.6, "Fractional Seconds in Time Values".

For `TIME` values specified as strings that include a time part delimiter, it is unnecessary to specify two digits for hours, minutes, or seconds values that are less than `10`. `'8:3:2'` is the same as `'08:03:02'`.

9.1.4 Hexadecimal Literals

Hexadecimal literal values are written using `X'val'` or `0xval` notation, where *`val`* contains hexadecimal digits (`0..9`, `A..F`). Lettercase of the digits and of any leading `X` does not matter. A leading `0x` is case sensitive and cannot be written as `0X`.

Legal hexadecimal literals:

```
X'01AF'
X'01af'
x'01AF'
x'01af'
0x01AF
0x01af
```

Illegal hexadecimal literals:

```
X'0G'   (G is not a hexadecimal digit)
0X01AF  (0X must be written as 0x)
```

Values written using `X'val'` notation must contain an even number of digits or a syntax error occurs. To correct the problem, pad the value with a leading zero:

```
mysql> SET @s = X'FFF';
ERROR 1064 (42000): You have an error in your SQL syntax;
check the manual that corresponds to your MySQL server
```

```
version for the right syntax to use near 'X'FFF''

mysql> SET @s = X'0FFF';
Query OK, 0 rows affected (0.00 sec)
```

Values written using `0xval` notation that contain an odd number of digits are treated as having an extra leading `0`. For example, `0xaaa` is interpreted as `0x0aaa`.

By default, a hexadecimal literal is a binary string, where each pair of hexadecimal digits represents a character:

```
mysql> SELECT X'4D7953514C', CHARSET(X'4D7953514C');
+---------------+------------------------+
| X'4D7953514C' | CHARSET(X'4D7953514C') |
+---------------+------------------------+
| MySQL         | binary                 |
+---------------+------------------------+
mysql> SELECT 0x5461626c65, CHARSET(0x5461626c65);
+--------------+-----------------------+
| 0x5461626c65 | CHARSET(0x5461626c65) |
+--------------+-----------------------+
| Table        | binary                |
+--------------+-----------------------+
```

A hexadecimal literal may have an optional character set introducer and `COLLATE` clause, to designate it as a string that uses a particular character set and collation:

```
[_charset_name] X'val' [COLLATE collation_name]
```

Examples:

```
SELECT _latin1 X'4D7953514C';
SELECT _utf8 0x4D7953514C COLLATE utf8_danish_ci;
```

The examples use `X'val'` notation, but `0xval` notation permits introducers as well. For information about introducers, see Section 10.1.3.8, “Character Set Introducers”.

In numeric contexts, MySQL treats a hexadecimal literal like a `BIGINT` (64-bit integer). To ensure numeric treatment of a hexadecimal literal, use it in numeric context. Ways to do this include adding 0 or using `CAST(... AS UNSIGNED)`. For example, a hexadecimal literal assigned to a user-defined variable is a binary string by default. To assign the value as a number, use it in numeric context:

```
mysql> SET @v1 = X'41';
mysql> SET @v2 = X'41'+0;
mysql> SET @v3 = CAST(X'41' AS UNSIGNED);
mysql> SELECT @v1, @v2, @v3;
+------+------+------+
| @v1  | @v2  | @v3  |
+------+------+------+
| A    |   65 |   65 |
+------+------+------+
```

An empty hexadecimal value (`X''`) evaluates to a zero-length binary string. Converted to a number, it produces 0:

```
mysql> SELECT CHARSET(X''), LENGTH(X'');
+--------------+-------------+
| CHARSET(X'') | LENGTH(X'') |
+--------------+-------------+
| binary       |           0 |
+--------------+-------------+
mysql> SELECT X''+0;
+-------+
| X''+0 |
+-------+
```

```
|     0 |
+-------+
```

The `X'val'` notation is based on standard SQL. The `0x` notation is based on ODBC, for which hexadecimal strings are often used to supply values for `BLOB` columns.

To convert a string or a number to a string in hexadecimal format, use the `HEX()` function:

```
mysql> SELECT HEX('cat');
+------------+
| HEX('cat') |
+------------+
| 636174     |
+------------+
mysql> SELECT X'636174';
+-----------+
| X'636174' |
+-----------+
| cat       |
+-----------+
```

For hexadecimal literals, bit operations are considered numeric context, but bit operations permit numeric or binary string arguments in MySQL 8.0 and higher. To explicitly specify binary string context for hexadecimal literals, use a `_binary` introducer for at least one of the arguments:

```
mysql> SET @v1 = X'000D' | X'0BC0';
mysql> SET @v2 = _binary X'000D' | X'0BC0';
mysql> SELECT HEX(@v1), HEX(@v2);
+----------+----------+
| HEX(@v1) | HEX(@v2) |
+----------+----------+
| BCD      | 0BCD     |
+----------+----------+
```

The displayed result appears similar for both bit operations, but the result without `_binary` is a `BIGINT` value, whereas the result with `_binary` is a binary string. Due to the difference in result types, the displayed values differ: High-order 0 digits are not displayed for the numeric result.

9.1.5 Bit-Value Literals

Bit-value literals are written using `b'val'` or `0bval` notation. *`val`* is a binary value written using zeros and ones. Lettercase of any leading `b` does not matter. A leading `0b` is case sensitive and cannot be written as `0B`.

Legal bit-value literals:

```
b'01'
B'01'
0b01
```

Illegal bit-value literals:

```
b'2'    (2 is not a binary digit)
0B01    (0B must be written as 0b)
```

By default, a bit-value literal is a binary string:

```
mysql> SELECT b'1000001', CHARSET(b'1000001');
+------------+----------------------+
| b'1000001' | CHARSET(b'1000001') |
+------------+----------------------+
| A          | binary               |
+------------+----------------------+
mysql> SELECT 0b1100001, CHARSET(0b1100001);
```

```
+-----------+--------------------+
| 0b1100001 | CHARSET(0b1100001) |
+-----------+--------------------+
| a         | binary             |
+-----------+--------------------+
```

A bit-value literal may have an optional character set introducer and `COLLATE` clause, to designate it as a string that uses a particular character set and collation:

```
[_charset_name] b'val' [COLLATE collation_name]
```

Examples:

```
SELECT _latin1 b'1000001';
SELECT _utf8 0b1000001 COLLATE utf8_danish_ci;
```

The examples use `b'val'` notation, but `0bval` notation permits introducers as well. For information about introducers, see Section 10.1.3.8, “Character Set Introducers”.

In numeric contexts, MySQL treats a bit literal like an integer. To ensure numeric treatment of a bit literal, use it in numeric context. Ways to do this include adding 0 or using `CAST(... AS UNSIGNED)`. For example, a bit literal assigned to a user-defined variable is a binary string by default. To assign the value as a number, use it in numeric context:

```
mysql> SET @v1 = b'1100001';
mysql> SET @v2 = b'1100001'+0;
mysql> SET @v3 = CAST(b'1100001' AS UNSIGNED);
mysql> SELECT @v1, @v2, @v3;
+------+------+------+
| @v1  | @v2  | @v3  |
+------+------+------+
| a    |   97 |   97 |
+------+------+------+
```

An empty bit value (`b''`) evaluates to a zero-length binary string. Converted to a number, it produces 0:

```
mysql> SELECT CHARSET(b''), LENGTH(b'');
+--------------+-------------+
| CHARSET(b'') | LENGTH(b'') |
+--------------+-------------+
| binary       |           0 |
+--------------+-------------+
mysql> SELECT b''+0;
+-------+
| b''+0 |
+-------+
|     0 |
+-------+
```

Bit-value notation is convenient for specifying values to be assigned to `BIT` columns:

```
mysql> CREATE TABLE t (b BIT(8));
mysql> INSERT INTO t SET b = b'11111111';
mysql> INSERT INTO t SET b = b'1010';
mysql> INSERT INTO t SET b = b'0101';
```

Bit values in result sets are returned as binary values, which may not display well. To convert a bit value to printable form, use it in numeric context or use a conversion function such as `BIN()` or `HEX()`. High-order 0 digits are not displayed in the converted value.

```
mysql> SELECT b+0, BIN(b), OCT(b), HEX(b) FROM t;
+------+----------+--------+--------+
```

```
| b+0  | BIN(b)   | OCT(b) | HEX(b) |
+------+----------+--------+--------+
|  255 | 11111111 | 377    | FF     |
|   10 | 1010     | 12     | A      |
|    5 | 101      | 5      | 5      |
+------+----------+--------+--------+
```

For bit literals, bit operations are considered numeric context, but bit operations permit numeric or binary string arguments in MySQL 8.0 and higher. To explicitly specify binary string context for bit literals, use a `_binary` introducer for at least one of the arguments:

```
mysql> SET @v1 = b'000010101' | b'000101010';
mysql> SET @v2 = _binary b'000010101' | _binary b'000101010';
mysql> SELECT HEX(@v1), HEX(@v2);
+----------+----------+
| HEX(@v1) | HEX(@v2) |
+----------+----------+
| 3F       | 003F     |
+----------+----------+
```

The displayed result appears similar for both bit operations, but the result without `_binary` is a `BIGINT` value, whereas the result with `_binary` is a binary string. Due to the difference in result types, the displayed values differ: High-order 0 digits are not displayed for the numeric result.

9.1.6 Boolean Literals

The constants `TRUE` and `FALSE` evaluate to `1` and `0`, respectively. The constant names can be written in any lettercase.

```
mysql> SELECT TRUE, true, FALSE, false;
        -> 1, 1, 0, 0
```

9.1.7 NULL Values

The `NULL` value means “no data.” `NULL` can be written in any lettercase.

Be aware that the `NULL` value is different from values such as `0` for numeric types or the empty string for string types. For more information, see Section B.5.4.3, “Problems with NULL Values”.

For text file import or export operations performed with `LOAD DATA INFILE` or `SELECT ... INTO OUTFILE`, `NULL` is represented by the `\N` sequence. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

9.2 Schema Object Names

Certain objects within MySQL, including database, table, index, column, alias, view, stored procedure, partition, tablespace, and other object names are known as identifiers. This section describes the permissible syntax for identifiers in MySQL. Section 9.2.2, “Identifier Case Sensitivity”, describes which types of identifiers are case sensitive and under what conditions.

An identifier may be quoted or unquoted. If an identifier contains special characters or is a reserved word, you *must* quote it whenever you refer to it. (Exception: A reserved word that follows a period in a qualified name must be an identifier, so it need not be quoted.) Reserved words are listed at Section 9.3, “Keywords and Reserved Words”.

Identifiers are converted to Unicode internally. They may contain these characters:

- Permitted characters in unquoted identifiers:
  - ASCII: [0-9,a-z,A-Z$_] (basic Latin letters, digits 0-9, dollar, underscore)
  - Extended: U+0080 .. U+FFFF

- Permitted characters in quoted identifiers include the full Unicode Basic Multilingual Plane (BMP), except U+0000:
  - ASCII: U+0001 .. U+007F
  - Extended: U+0080 .. U+FFFF
- ASCII NUL (U+0000) and supplementary characters (U+10000 and higher) are not permitted in quoted or unquoted identifiers.
- Identifiers may begin with a digit but unless quoted may not consist solely of digits.
- Database, table, and column names cannot end with space characters.

The identifier quote character is the backtick (`):

```
mysql> SELECT * FROM `select` WHERE `select`.id > 100;
```

If the ANSI_QUOTES SQL mode is enabled, it is also permissible to quote identifiers within double quotation marks:

```
mysql> CREATE TABLE "test" (col INT);
ERROR 1064: You have an error in your SQL syntax...
mysql> SET sql_mode='ANSI_QUOTES';
mysql> CREATE TABLE "test" (col INT);
Query OK, 0 rows affected (0.00 sec)
```

The ANSI_QUOTES mode causes the server to interpret double-quoted strings as identifiers. Consequently, when this mode is enabled, string literals must be enclosed within single quotation marks. They cannot be enclosed within double quotation marks. The server SQL mode is controlled as described in Section 5.1.8, “Server SQL Modes”.

Identifier quote characters can be included within an identifier if you quote the identifier. If the character to be included within the identifier is the same as that used to quote the identifier itself, then you need to double the character. The following statement creates a table named a`b that contains a column named c"d:

```
mysql> CREATE TABLE `a``b` (`c"d` INT);
```

In the select list of a query, a quoted column alias can be specified using identifier or string quoting characters:

```
mysql> SELECT 1 AS `one`, 2 AS 'two';
+-----+-----+
| one | two |
+-----+-----+
|   1 |   2 |
+-----+-----+
```

Elsewhere in the statement, quoted references to the alias must use identifier quoting or the reference is treated as a string literal.

It is recommended that you do not use names that begin with *Me* or *MeN*, where *M* and *N* are integers. For example, avoid using 1e as an identifier, because an expression such as 1e+3 is ambiguous. Depending on context, it might be interpreted as the expression 1e + 3 or as the number 1e+3.

Be careful when using MD5() to produce table names because it can produce names in illegal or ambiguous formats such as those just described.

A user variable cannot be used directly in an SQL statement as an identifier or as part of an identifier. See Section 9.4, “User-Defined Variables”, for more information and examples of workarounds.

Special characters in database and table names are encoded in the corresponding file system names as described in Section 9.2.3, “Mapping of Identifiers to File Names”.

The following table describes the maximum length for each type of identifier.

Identifier	Maximum Length (characters)
Database	64
Table	64
Column	64
Index	64
Constraint	64
Stored Program	64
View	64
Tablespace	64
Server	64
Log File Group	64
Alias	256 (see exception following table)
Compound Statement Label	16
User-Defined Variable	64

Aliases for column names in `CREATE VIEW` statements are checked against the maximum column length of 64 characters (not the maximum alias length of 256 characters).

Identifiers are stored using Unicode (UTF-8). This applies to identifiers in table definitions and to identifiers stored in the grant tables in the `mysql` database. The sizes of the identifier string columns in the grant tables are measured in characters. You can use multibyte characters without reducing the number of characters permitted for values stored in these columns. As indicated earlier, the permissible Unicode characters are those in the Basic Multilingual Plane (BMP). Supplementary characters are not permitted.

9.2.1 Identifier Qualifiers

Object names may be unqualified or qualified. An unqualified name is permitted in contexts where interpretation of the name is unambiguous. A qualified name includes at least one qualifier to clarify the interpretive context by overriding a default context or providing missing context.

For example, this statement creates a table using the unqualified name `t1`:

```
CREATE TABLE t1 (i INT);
```

Because `t1` includes no qualifier to specify a database, the statement creates the table in the default database. If there is no default database, an error occurs.

This statement creates a table using the qualified name `db1.t1`:

```
CREATE TABLE db1.t1 (i INT);
```

Because `db1.t1` includes a database qualifier `db1`, the statement creates `t1` in the database named `db1`, regardless of the default database. The qualifier *must* be specified if there is no default database. The qualifier *may* be specified if there is a default database, to specify a database different from the default, or to make the database explicit if the default is the same as the one specified.

Qualifiers have these characteristics:

- An unqualified name consists of a single identifier. A qualified name consists of multiple identifiers.

- The components of a multiple-part name must be separated by period (`.`) characters. The initial parts of a multiple-part name act as qualifiers that affect the context within which to interpret the final identifier.
- The qualifier character is a separate token and need not be contiguous with the associated identifiers. For example, *`tbl_name.col_name`* and *`tbl_name . col_name`* are equivalent.
- If any components of a multiple-part name require quoting, quote them individually rather than quoting the name as a whole. For example, write `` `my-table`.`my-column` ``, not `` `my-table.my-column` ``.
- A reserved word that follows a period in a qualified name must be an identifier, so in that context it need not be quoted.

The permitted qualifiers for object names depend on the object type:

- A database name is fully qualified and takes no qualifier:

```
CREATE DATABASE db1;
```

- A table, view, or stored program name may be given a database-name qualifier. Examples of unqualified and qualified names in `CREATE` statements:

```
CREATE TABLE mytable ...;
CREATE VIEW myview ...;
CREATE PROCEDURE myproc ...;
CREATE FUNCTION myfunc ...;
CREATE EVENT myevent ...;

CREATE TABLE mydb.mytable ...;
CREATE VIEW mydb.myview ...;
CREATE PROCEDURE mydb.myproc ...;
CREATE FUNCTION mydb.myfunc ...;
CREATE EVENT mydb.myevent ...;
```

- A trigger is associated with a table, so any qualifier applies to the table name:

```
CREATE TRIGGER mytrigger ... ON mytable ...;

CREATE TRIGGER mytrigger ... ON mydb.mytable ...;
```

- A column name may be given multiple qualifiers to indicate context in statements that reference it, as shown in the following table.

Column Reference	Meaning
`col_name`	Column *`col_name`* from whichever table used in the statement contains a column of that name
`tbl_name.col_name`	Column *`col_name`* from table *`tbl_name`* of the default database
`db_name.tbl_name.col_name`	Column *`col_name`* from table *`tbl_name`* of the database *`db_name`*

In other words, a column name may be given a table-name qualifier, which itself may be given a database-name qualifier. Examples of unqualified and qualified column references in `SELECT` statements:

```
SELECT c1 FROM mytable
WHERE c2 > 100;

SELECT mytable.c1 FROM mytable
```

```
WHERE mytable.c2 > 100;

SELECT mydb.mytable.c1 FROM mydb.mytable
WHERE mydb.mytable.c2 > 100;
```

You need not specify a qualifier for an object reference in a statement unless the unqualified reference is ambiguous. Suppose that column `c1` occurs only in table `t1`, `c2` only in `t2`, and `c` in both `t1` and `t2`. Any unqualified reference to `c` is ambiguous in a statement that refers to both tables and must be qualified as `t1.c` or `t2.c` to indicate which table you mean:

```
SELECT c1, c2, t1.c FROM t1 INNER JOIN t2
WHERE t2.c > 100;
```

Similarly, to retrieve from a table `t` in database `db1` and from a table `t` in database `db2` in the same statement, you must qualify the table references: For references to columns in those tables, qualifiers are required only for column names that appear in both tables. Suppose that column `c1` occurs only in table `db1.t`, `c2` only in `db2.t`, and `c` in both `db1.t` and `db2.t`. In this case, `c` is ambiguous and must be qualified but `c1` and `c2` need not be:

```
SELECT c1, c2, db1.t.c FROM db1.t INNER JOIN db2.t
WHERE db2.t.c > 100;
```

Table aliases enable qualified column references to be written more simply:

```
SELECT c1, c2, t1.c FROM db1.t AS t1 INNER JOIN db2.t AS t2
WHERE t2.c > 100;
```

9.2.2 Identifier Case Sensitivity

In MySQL, databases correspond to directories within the data directory. Each table within a database corresponds to at least one file within the database directory (and possibly more, depending on the storage engine). Triggers also correspond to files. Consequently, the case sensitivity of the underlying operating system plays a part in the case sensitivity of database, table, and trigger names. This means such names are not case sensitive in Windows, but are case sensitive in most varieties of Unix. One notable exception is OS X, which is Unix-based but uses a default file system type (HFS+) that is not case sensitive. However, OS X also supports UFS volumes, which are case sensitive just as on any Unix. See Section 1.8.1, “MySQL Extensions to Standard SQL”. The `lower_case_table_names` system variable also affects how the server handles identifier case sensitivity, as described later in this section.

Note

Although database, table, and trigger names are not case sensitive on some platforms, you should not refer to one of these using different cases within the same statement. The following statement would not work because it refers to a table both as `my_table` and as `MY_TABLE`:

```
mysql> SELECT * FROM my_table WHERE MY_TABLE.col=1;
```

Column, index, stored routine, and event names are not case sensitive on any platform, nor are column aliases.

However, names of logfile groups are case sensitive. This differs from standard SQL.

By default, table aliases are case sensitive on Unix, but not so on Windows or OS X. The following statement would not work on Unix, because it refers to the alias both as `a` and as `A`:

```
mysql> SELECT col_name FROM tbl_name AS a
    -> WHERE a.col_name = 1 OR A.col_name = 2;
```

However, this same statement is permitted on Windows. To avoid problems caused by such differences, it is best to adopt a consistent convention, such as always creating and referring to databases and tables using lowercase names. This convention is recommended for maximum portability and ease of use.

How table and database names are stored on disk and used in MySQL is affected by the `lower_case_table_names` system variable, which you can set when starting `mysqld`. `lower_case_table_names` can take the values shown in the following table. This variable does *not* affect case sensitivity of trigger identifiers. On Unix, the default value of `lower_case_table_names` is 0. On Windows, the default value is 1. On OS X, the default value is 2.

Value	Meaning
`0`	Table and database names are stored on disk using the lettercase specified in the `CREATE TABLE` or `CREATE DATABASE` statement. Name comparisons are case sensitive. You should *not* set this variable to 0 if you are running MySQL on a system that has case-insensitive file names (such as Windows or OS X). If you force this variable to 0 with `--lower-case-table-names=0` on a case-insensitive file system and access `MyISAM` tablenames using different lettercases, index corruption may result.
`1`	Table names are stored in lowercase on disk and name comparisons are not case sensitive. MySQL converts all table names to lowercase on storage and lookup. This behavior also applies to database names and table aliases.
`2`	Table and database names are stored on disk using the lettercase specified in the `CREATE TABLE` or `CREATE DATABASE` statement, but MySQL converts them to lowercase on lookup. Name comparisons are not case sensitive. This works *only* on file systems that are not case sensitive! `InnoDB` table names are stored in lowercase, as for `lower_case_table_names=1`.

If you are using MySQL on only one platform, you do not normally have to change the `lower_case_table_names` variable from its default value. However, you may encounter difficulties if you want to transfer tables between platforms that differ in file system case sensitivity. For example, on Unix, you can have two different tables named `my_table` and `MY_TABLE`, but on Windows these two names are considered identical. To avoid data transfer problems arising from lettercase of database or table names, you have two options:

- Use `lower_case_table_names=1` on all systems. The main disadvantage with this is that when you use `SHOW TABLES` or `SHOW DATABASES`, you do not see the names in their original lettercase.
- Use `lower_case_table_names=0` on Unix and `lower_case_table_names=2` on Windows. This preserves the lettercase of database and table names. The disadvantage of this is that you must ensure that your statements always refer to your database and table names with the correct lettercase on Windows. If you transfer your statements to Unix, where lettercase is significant, they do not work if the lettercase is incorrect.

  Exception: If you are using `InnoDB` tables and you are trying to avoid these data transfer problems, you should set `lower_case_table_names` to 1 on all platforms to force names to be converted to lowercase.

If you plan to set the `lower_case_table_names` system variable to 1 on Unix, you must first convert your old database and table names to lowercase before stopping `mysqld` and restarting it with the new variable setting. To do this for an individual table, use `RENAME TABLE`:

```
RENAME TABLE T1 TO t1;
```

To convert one or more entire databases, dump them before setting `lower_case_table_names`, then drop the databases, and reload them after setting `lower_case_table_names`:

1. Use `mysqldump` to dump each database:

   ```
   mysqldump --databases db1 > db1.sql
   ```

```
mysqldump --databases db2 > db2.sql
...
```

Do this for each database that must be recreated.

2. Use `DROP DATABASE` to drop each database.
3. Stop the server, set `lower_case_table_names`, and restart the server.
4. Reload the dump file for each database. Because `lower_case_table_names` is set, each database and table name will be converted to lowercase as it is recreated:

```
mysql < db1.sql
mysql < db2.sql
...
```

Object names may be considered duplicates if their uppercase forms are equal according to a binary collation. That is true for names of cursors, conditions, procedures, functions, savepoints, stored routine parameters, stored program local variables, and plugins. It is not true for names of columns, constraints, databases, partitions, statements prepared with `PREPARE`, tables, triggers, users, and user-defined variables.

File system case sensitivity can affect searches in string columns of `INFORMATION_SCHEMA` tables. For more information, see Section 10.1.8.7, “Using Collation in INFORMATION_SCHEMA Searches”.

9.2.3 Mapping of Identifiers to File Names

There is a correspondence between database and table identifiers and names in the file system. For the basic structure, MySQL represents each database as a directory in the data directory, and depending upon the storage engine, each table may be represented by one or more files in the appropriate database directory.

For the data and index files, the exact representation on disk is storage engine specific. These files may be stored in the database directory, or the information may be stored in a separate file. `InnoDB` data is stored in the InnoDB data files. If you are using tablespaces with `InnoDB`, then the specific tablespace files you create are used instead.

Any character is legal in database or table identifiers except ASCII NUL (`X'00'`). MySQL encodes any characters that are problematic in the corresponding file system objects when it creates database directories or table files:

- Basic Latin letters (`a..zA..Z`), digits (`0..9`) and underscore (_) are encoded as is. Consequently, their case sensitivity directly depends on file system features.
- All other national letters from alphabets that have uppercase/lowercase mapping are encoded as shown in the following table. Values in the Code Range column are UCS-2 values.

Code Range	Pattern	Number	Used	Unused	Blocks
00C0..017F	[@][0..4][g..z]	5*20= 100	97	3	Latin-1 Supplement + Latin Extended-A
0370..03FF	[@][5..9][g..z]	5*20= 100	88	12	Greek and Coptic
0400..052F	[@][g..z][0..6]	20*7= 140	137	3	Cyrillic + Cyrillic Supplement
0530..058F	[@][g..z][7..8]	20*2= 40	38	2	Armenian
2160..217F	[@][g..z][9]	20*1= 20	16	4	Number Forms
0180..02AF	[@][g..z][a..k]	20*11=220	203	17	Latin Extended-B + IPA Extensions
1E00..1EFF	[@][g..z][l..r]	20*7= 140	136	4	Latin Extended Additional

Code Range	Pattern	Number	Used	Unused	Blocks
1F00..1FFF	[@][g..z][s..z]	20*8= 160	144	16	Greek Extended
....	[@][a..f][g..z]	6*20= 120	0	120	RESERVED
24B6..24E9	[@][@][a..z]	26	26	0	Enclosed Alphanumerics
FF21..FF5A	[@][a..z][@]	26	26	0	Halfwidth and Fullwidth forms

One of the bytes in the sequence encodes lettercase. For example: `LATIN CAPITAL LETTER A WITH GRAVE` is encoded as `@0G`, whereas `LATIN SMALL LETTER A WITH GRAVE` is encoded as `@0g`. Here the third byte (`G` or `g`) indicates lettercase. (On a case-insensitive file system, both letters will be treated as the same.)

For some blocks, such as Cyrillic, the second byte determines lettercase. For other blocks, such as Latin1 Supplement, the third byte determines lettercase. If two bytes in the sequence are letters (as in Greek Extended), the leftmost letter character stands for lettercase. All other letter bytes must be in lowercase.

- All nonletter characters except underscore (_), as well as letters from alphabets that do not have uppercase/lowercase mapping (such as Hebrew) are encoded using hexadecimal representation using lowercase letters for hexadecimal digits `a..f`:

```
0x003F -> @003f
0xFFFF -> @ffff
```

The hexadecimal values correspond to character values in the `ucs2` double-byte character set.

On Windows, some names such as `nul`, `prn`, and `aux` are encoded by appending `@@@` to the name when the server creates the corresponding file or directory. This occurs on all platforms for portability of the corresponding database object between platforms.

9.2.4 Function Name Parsing and Resolution

MySQL 8.0 supports built-in (native) functions, user-defined functions (UDFs), and stored functions. This section describes how the server recognizes whether the name of a built-in function is used as a function call or as an identifier, and how the server determines which function to use in cases when functions of different types exist with a given name.

Built-In Function Name Parsing

The parser uses default rules for parsing names of built-in functions. These rules can be changed by enabling the `IGNORE_SPACE` SQL mode.

When the parser encounters a word that is the name of a built-in function, it must determine whether the name signifies a function call or is instead a nonexpression reference to an identifier such as a table or column name. For example, in the following statements, the first reference to `count` is a function call, whereas the second reference is a table name:

```
SELECT COUNT(*) FROM mytable;
CREATE TABLE count (i INT);
```

The parser should recognize the name of a built-in function as indicating a function call only when parsing what is expected to be an expression. That is, in nonexpression context, function names are permitted as identifiers.

However, some built-in functions have special parsing or implementation considerations, so the parser uses the following rules by default to distinguish whether their names are being used as function calls or as identifiers in nonexpression context:

- To use the name as a function call in an expression, there must be no whitespace between the name and the following `(` parenthesis character.
- Conversely, to use the function name as an identifier, it must not be followed immediately by a parenthesis.

The requirement that function calls be written with no whitespace between the name and the parenthesis applies only to the built-in functions that have special considerations. `COUNT` is one such name. The `sql/lex.h` source file lists the names of these special functions for which following whitespace determines their interpretation: names defined by the `SYM_FN()` macro in the `symbols[]` array.

In MySQL 8.0, there are about 30 such function names. You may find it easiest to treat the no-whitespace requirement as applying to all function calls.

The following table names the functions that are affected by the `IGNORE_SPACE` setting and listed as special in the `sql/lex.h` source file.

ADDDATE	BIT_AND	BIT_OR	BIT_XOR
CAST	COUNT	CURDATE	CURTIME
DATE_ADD	DATE_SUB	EXTRACT	GROUP_CONCAT
MAX	MID	MIN	NOW
POSITION	SESSION_USER	STD	STDDEV
STDDEV_POP	STDDEV_SAMP	SUBDATE	SUBSTR
SUBSTRING	SUM	SYSDATE	SYSTEM_USER
TRIM	VARIANCE	VAR_POP	VAR_SAMP

For functions not listed as special in `sql/lex.h`, whitespace does not matter. They are interpreted as function calls only when used in expression context and may be used freely as identifiers otherwise. `ASCII` is one such name. However, for these nonaffected function names, interpretation may vary in expression context: *`func_name ()`* is interpreted as a built-in function if there is one with the given name; if not, *`func_name ()`* is interpreted as a user-defined function or stored function if one exists with that name.

The `IGNORE_SPACE` SQL mode can be used to modify how the parser treats function names that are whitespace-sensitive:

- With `IGNORE_SPACE` disabled, the parser interprets the name as a function call when there is no whitespace between the name and the following parenthesis. This occurs even when the function name is used in nonexpression context:

  ```
  mysql> CREATE TABLE count(i INT);
  ERROR 1064 (42000): You have an error in your SQL syntax ...
  near 'count(i INT)'
  ```

  To eliminate the error and cause the name to be treated as an identifier, either use whitespace following the name or write it as a quoted identifier (or both):

  ```
  CREATE TABLE count (i INT);
  CREATE TABLE `count`(i INT);
  CREATE TABLE `count` (i INT);
  ```

- With `IGNORE_SPACE` enabled, the parser loosens the requirement that there be no whitespace between the function name and the following parenthesis. This provides more flexibility in writing function calls. For example, either of the following function calls are legal:

  ```
  ```

```
SELECT COUNT(*) FROM mytable;
SELECT COUNT (*) FROM mytable;
```

However, enabling `IGNORE_SPACE` also has the side effect that the parser treats the affected function names as reserved words (see Section 9.3, “Keywords and Reserved Words”). This means that a space following the name no longer signifies its use as an identifier. The name can be used in function calls with or without following whitespace, but causes a syntax error in nonexpression context unless it is quoted. For example, with `IGNORE_SPACE` enabled, both of the following statements fail with a syntax error because the parser interprets `count` as a reserved word:

```
CREATE TABLE count(i INT);
CREATE TABLE count (i INT);
```

To use the function name in nonexpression context, write it as a quoted identifier:

```
CREATE TABLE `count`(i INT);
CREATE TABLE `count` (i INT);
```

To enable the `IGNORE_SPACE` SQL mode, use this statement:

```
SET sql_mode = 'IGNORE_SPACE';
```

`IGNORE_SPACE` is also enabled by certain other composite modes such as `ANSI` that include it in their value:

```
SET sql_mode = 'ANSI';
```

Check Section 5.1.8, “Server SQL Modes”, to see which composite modes enable `IGNORE_SPACE`.

To minimize the dependency of SQL code on the `IGNORE_SPACE` setting, use these guidelines:

- Avoid creating UDFs or stored functions that have the same name as a built-in function.
- Avoid using function names in nonexpression context. For example, these statements use `count` (one of the affected function names affected by `IGNORE_SPACE`), so they fail with or without whitespace following the name if `IGNORE_SPACE` is enabled:

  ```
  CREATE TABLE count(i INT);
  CREATE TABLE count (i INT);
  ```

  If you must use a function name in nonexpression context, write it as a quoted identifier:

  ```
  CREATE TABLE `count`(i INT);
  CREATE TABLE `count` (i INT);
  ```

Function Name Resolution

The following rules describe how the server resolves references to function names for function creation and invocation:

- Built-in functions and user-defined functions

  An error occurs if you try to create a UDF with the same name as a built-in function.

- Built-in functions and stored functions

  It is possible to create a stored function with the same name as a built-in function, but to invoke the stored function it is necessary to qualify it with a schema name. For example, if you create a stored function named `PI` in the `test` schema, invoke it as `test.PI()` because the server resolves

`PI()` without a qualifier as a reference to the built-in function. The server generates a warning if the stored function name collides with a built-in function name. The warning can be displayed with `SHOW WARNINGS`.

- User-defined functions and stored functions

  User-defined functions and stored functions share the same namespace, so you cannot create a UDF and a stored function with the same name.

The preceding function name resolution rules have implications for upgrading to versions of MySQL that implement new built-in functions:

- If you have already created a user-defined function with a given name and upgrade MySQL to a version that implements a new built-in function with the same name, the UDF becomes inaccessible. To correct this, use `DROP FUNCTION` to drop the UDF and `CREATE FUNCTION` to re-create the UDF with a different nonconflicting name. Then modify any affected code to use the new name.
- If a new version of MySQL implements a built-in function with the same name as an existing stored function, you have two choices: Rename the stored function to use a nonconflicting name, or change calls to the function so that they use a schema qualifier (that is, use *`schema_name.func_name()`* syntax). In either case, modify any affected code accordingly.

9.3 Keywords and Reserved Words

Keywords are words that have significance in SQL. Certain keywords, such as `SELECT`, `DELETE`, or `BIGINT`, are reserved and require special treatment for use as identifiers such as table and column names. This may also be true for the names of built-in functions.

Nonreserved keywords are permitted as identifiers without quoting. Reserved words are permitted as identifiers if you quote them as described in Section 9.2, “Schema Object Names”:

```
mysql> CREATE TABLE interval (begin INT, end INT);
ERROR 1064 (42000): You have an error in your SQL syntax ...
near 'interval (begin INT, end INT)'
```

`BEGIN` and `END` are keywords but not reserved, so their use as identifiers does not require quoting. `INTERVAL` is a reserved keyword and must be quoted to be used as an identifier:

```
mysql> CREATE TABLE `interval` (begin INT, end INT);
Query OK, 0 rows affected (0.01 sec)
```

Exception: A word that follows a period in a qualified name must be an identifier, so it need not be quoted even if it is reserved:

```
mysql> CREATE TABLE mydb.interval (begin INT, end INT);
Query OK, 0 rows affected (0.01 sec)
```

Names of built-in functions are permitted as identifiers but may require care to be used as such. For example, `COUNT` is acceptable as a column name. However, by default, no whitespace is permitted in function invocations between the function name and the following `(` character. This requirement enables the parser to distinguish whether the name is used in a function call or in nonfunction context. For further details on recognition of function names, see Section 9.2.4, “Function Name Parsing and Resolution”.

The following table shows the keywords and reserved words in MySQL 8.0, along with changes to individual words from version to version. Reserved keywords are marked with (R). In addition, `_FILENAME` is reserved.

At some point, you might upgrade to a higher version, so it is a good idea to have a look at future reserved words, too. You can find these in the manuals that cover higher versions of MySQL. Most of

the reserved words in the table are forbidden by standard SQL as column or table names (for example, GROUP). A few are reserved because MySQL needs them and uses a yacc parser.

Table 9.2 Keywords and Reserved Words in MySQL 8.0

ACCESSIBLE (R)	ACCOUNT	ACTION
ADD (R)	ADMIN (R)	AFTER
AGAINST	AGGREGATE	ALGORITHM
ALL (R)	ALTER (R)	ALWAYS
ANALYSE[a]	ANALYZE (R)	AND (R)
ANY	AS (R)	ASC (R)
ASCII	ASENSITIVE (R)	AT
AUTOEXTEND_SIZE	AUTO_INCREMENT	AVG
AVG_ROW_LENGTH	BACKUP	BEFORE (R)
BEGIN	BETWEEN (R)	BIGINT (R)
BINARY (R)	BINLOG	BIT
BLOB (R)	BLOCK	BOOL
BOOLEAN	BOTH (R)	BTREE
BUCKETS[b]	BY (R)	BYTE
CACHE	CALL (R)	CASCADE (R)
CASCADED	CASE (R)	CATALOG_NAME
CHAIN	CHANGE (R)	CHANGED
CHANNEL	CHAR (R)	CHARACTER (R)
CHARSET	CHECK (R)	CHECKSUM
CIPHER	CLASS_ORIGIN	CLIENT
CLONE[c]	CLOSE	COALESCE
CODE	COLLATE (R)	COLLATION
COLUMN (R)	COLUMNS	COLUMN_FORMAT
COLUMN_NAME	COMMENT	COMMIT
COMMITTED	COMPACT	COMPLETION
COMPONENT	COMPRESSED	COMPRESSION
CONCURRENT	CONDITION (R)	CONNECTION
CONSISTENT	CONSTRAINT (R)	CONSTRAINT_CATALOG
CONSTRAINT_NAME	CONSTRAINT_SCHEMA	CONTAINS
CONTEXT	CONTINUE (R)	CONVERT (R)
CPU	CREATE (R)	CROSS (R)
CUBE[d] (R)	CUME_DIST[e] (R)	CURRENT
CURRENT_DATE (R)	CURRENT_TIME (R)	CURRENT_TIMESTAMP (R)
CURRENT_USER (R)	CURSOR (R)	CURSOR_NAME
DATA	DATABASE (R)	DATABASES (R)
DATAFILE	DATE	DATETIME
DAY	DAY_HOUR (R)	DAY_MICROSECOND (R)
DAY_MINUTE (R)	DAY_SECOND (R)	DEALLOCATE
DEC (R)	DECIMAL (R)	DECLARE (R)

DEFAULT (R)	DEFAULT_AUTH	DEFINER
DELAYED (R)	DELAY_KEY_WRITE	DELETE (R)
DENSE_RANK[f] (R)	DESC (R)	DESCRIBE (R)
DES_KEY_FILE	DETERMINISTIC (R)	DIAGNOSTICS
DIRECTORY	DISABLE	DISCARD
DISK	DISTINCT (R)	DISTINCTROW (R)
DIV (R)	DO	DOUBLE (R)
DROP (R)	DUAL (R)	DUMPFILE
DUPLICATE	DYNAMIC	EACH (R)
ELSE (R)	ELSEIF (R)	ENABLE
ENCLOSED (R)	ENCRYPTION	END
ENDS	ENGINE	ENGINES
ENUM	ERROR	ERRORS
ESCAPE	ESCAPED (R)	EVENT
EVENTS	EVERY	EXCEPT (R)
EXCHANGE	EXCLUDE[g]	EXECUTE
EXISTS (R)	EXIT (R)	EXPANSION
EXPIRE	EXPLAIN (R)	EXPORT
EXTENDED	EXTENT_SIZE	FALSE (R)
FAST	FAULTS	FETCH (R)
FIELDS	FILE	FILE_BLOCK_SIZE
FILTER	FIRST	FIRST_VALUE[h] (R)
FIXED	FLOAT (R)	FLOAT4 (R)
FLOAT8 (R)	FLUSH	FOLLOWING[i]
FOLLOWS	FOR (R)	FORCE (R)
FOREIGN (R)	FORMAT	FOUND
FROM (R)	FULL	FULLTEXT (R)
FUNCTION[j] (R)	GENERAL	GENERATED (R)
GEOMETRY	GEOMETRYCOLLECTION	GET (R)
GET_FORMAT	GLOBAL	GRANT (R)
GRANTS	GROUP (R)	GROUPING[k] (R)
GROUPS[l] (R)	GROUP_REPLICATION	HANDLER
HASH	HAVING (R)	HELP
HIGH_PRIORITY (R)	HISTOGRAM[m]	HOST
HOSTS	HOUR	HOUR_MICROSECOND (R)
HOUR_MINUTE (R)	HOUR_SECOND (R)	IDENTIFIED
IF (R)	IGNORE (R)	IGNORE_SERVER_IDS
IMPORT	IN (R)	INDEX (R)
INDEXES	INFILE (R)	INITIAL_SIZE
INNER (R)	INOUT (R)	INSENSITIVE (R)
INSERT (R)	INSERT_METHOD	INSTALL
INSTANCE	INT (R)	INT1 (R)

INT2 (R)	INT3 (R)	INT4 (R)
INT8 (R)	INTEGER (R)	INTERVAL (R)
INTO (R)	INVISIBLE	INVOKER
IO	IO_AFTER_GTIDS (R)	IO_BEFORE_GTIDS (R)
IO_THREAD	IPC	IS (R)
ISOLATION	ISSUER	ITERATE (R)
JOIN (R)	JSON	KEY (R)
KEYS (R)	KEY_BLOCK_SIZE	KILL (R)
LAG[n] (R)	LANGUAGE	LAST
LAST_VALUE[o] (R)	LEAD[p] (R)	LEADING (R)
LEAVE (R)	LEAVES	LEFT (R)
LESS	LEVEL	LIKE (R)
LIMIT (R)	LINEAR (R)	LINES (R)
LINESTRING	LIST	LOAD (R)
LOCAL	LOCALTIME (R)	LOCALTIMESTAMP (R)
LOCK (R)	LOCKED[q]	LOCKS
LOGFILE	LOGS	LONG (R)
LONGBLOB (R)	LONGTEXT (R)	LOOP (R)
LOW_PRIORITY (R)	MASTER	MASTER_AUTO_POSITION
MASTER_BIND (R)	MASTER_CONNECT_RETRY	MASTER_DELAY
MASTER_HEARTBEAT_PERIOD	MASTER_HOST	MASTER_LOG_FILE
MASTER_LOG_POS	MASTER_PASSWORD	MASTER_PORT
MASTER_RETRY_COUNT	MASTER_SERVER_ID	MASTER_SSL
MASTER_SSL_CA	MASTER_SSL_CAPATH	MASTER_SSL_CERT
MASTER_SSL_CIPHER	MASTER_SSL_CRL	MASTER_SSL_CRLPATH
MASTER_SSL_KEY	MASTER_SSL_VERIFY_SERVER_CERT (R)	MASTER_TLS_VERSION
MASTER_USER	MATCH (R)	MAXVALUE (R)
MAX_CONNECTIONS_PER_HOUR	MAX_QUERIES_PER_HOUR	MAX_ROWS
MAX_SIZE	MAX_UPDATES_PER_HOUR	MAX_USER_CONNECTIONS
MEDIUM	MEDIUMBLOB (R)	MEDIUMINT (R)
MEDIUMTEXT (R)	MEMORY	MERGE
MESSAGE_TEXT	MICROSECOND	MIDDLEINT (R)
MIGRATE	MINUTE	MINUTE_MICROSECOND (R)
MINUTE_SECOND (R)	MIN_ROWS	MOD (R)
MODE	MODIFIES (R)	MODIFY
MONTH	MULTILINESTRING	MULTIPOINT
MULTIPOLYGON	MUTEX	MYSQL_ERRNO
NAME	NAMES	NATIONAL
NATURAL (R)	NCHAR	NDB
NDBCLUSTER	NEVER	NEW
NEXT	NO	NODEGROUP

NONE	NOT (R)	NOWAIT[r]
NO_WAIT	NO_WRITE_TO_BINLOG (R)	NTH_VALUE[s] (R)
NTILE[t] (R)	NULL (R)	NULLS[u]
NUMBER	NUMERIC (R)	NVARCHAR
OF[v] (R)	OFFSET	ON (R)
ONE	ONLY	OPEN
OPTIMIZE (R)	OPTIMIZER_COSTS (R)	OPTION (R)
OPTIONALLY (R)	OPTIONS	OR (R)
ORDER (R)	OTHERS[w]	OUT (R)
OUTER (R)	OUTFILE (R)	OVER[x] (R)
OWNER	PACK_KEYS	PAGE
PARSER	PARTIAL	PARTITION (R)
PARTITIONING	PARTITIONS	PASSWORD
PERCENT_RANK[y] (R)	PERSIST (R)	PERSIST_ONLY[z] (R)
PHASE	PLUGIN	PLUGINS
PLUGIN_DIR	POINT	POLYGON
PORT	PRECEDES	PRECEDING[aa]
PRECISION (R)	PREPARE	PRESERVE
PREV	PRIMARY (R)	PRIVILEGES
PROCEDURE (R)	PROCESSLIST	PROFILE
PROFILES	PROXY	PURGE (R)
QUARTER	QUERY	QUICK
RANGE (R)	RANK[ab] (R)	READ (R)
READS (R)	READ_ONLY	READ_WRITE (R)
REAL (R)	REBUILD	RECOVER
RECURSIVE[ac] (R)	REDOFILE	REDO_BUFFER_SIZE
REDUNDANT	REFERENCES (R)	REGEXP (R)
RELAY	RELAYLOG	RELAY_LOG_FILE
RELAY_LOG_POS	RELAY_THREAD	RELEASE (R)
RELOAD	REMOTE[ad]	REMOVE
RENAME (R)	REORGANIZE	REPAIR
REPEAT (R)	REPEATABLE	REPLACE (R)
REPLICATE_DO_DB	REPLICATE_DO_TABLE	REPLICATE_IGNORE_DB
REPLICATE_IGNORE_TABLE	REPLICATE_REWRITE_DB	REPLICATE_WILD_DO_TABLE
REPLICATE_WILD_IGNORE_TABLE	REPLICATION	REQUIRE (R)
RESET	RESIGNAL (R)	RESPECT[ae]
RESTORE	RESTRICT (R)	RESUME
RETURN (R)	RETURNED_SQLSTATE	RETURNS
REVERSE	REVOKE (R)	RIGHT (R)
RLIKE (R)	ROLE[af]	ROLLBACK
ROLLUP	ROTATE	ROUTINE
ROW[ag] (R)	ROWS[ah] (R)	ROW_COUNT

ROW_FORMAT	ROW_NUMBER[ai] (R)	RTREE
SAVEPOINT	SCHEDULE	SCHEMA (R)
SCHEMAS (R)	SCHEMA_NAME	SECOND
SECOND_MICROSECOND (R)	SECURITY	SELECT (R)
SENSITIVE (R)	SEPARATOR (R)	SERIAL
SERIALIZABLE	SERVER	SESSION
SET (R)	SHARE	SHOW (R)
SHUTDOWN	SIGNAL (R)	SIGNED
SIMPLE	SKIP[aj]	SLAVE
SLOW	SMALLINT (R)	SNAPSHOT
SOCKET	SOME	SONAME
SOUNDS	SOURCE	SPATIAL (R)
SPECIFIC (R)	SQL (R)	SQLEXCEPTION (R)
SQLSTATE (R)	SQLWARNING (R)	SQL_AFTER_GTIDS
SQL_AFTER_MTS_GAPS	SQL_BEFORE_GTIDS	SQL_BIG_RESULT (R)
SQL_BUFFER_RESULT	SQL_CACHE	SQL_CALC_FOUND_ROWS (R)
SQL_NO_CACHE	SQL_SMALL_RESULT (R)	SQL_THREAD
SQL_TSI_DAY	SQL_TSI_HOUR	SQL_TSI_MINUTE
SQL_TSI_MONTH	SQL_TSI_QUARTER	SQL_TSI_SECOND
SQL_TSI_WEEK	SQL_TSI_YEAR	SSL (R)
STACKED	START	STARTING (R)
STARTS	STATS_AUTO_RECALC	STATS_PERSISTENT
STATS_SAMPLE_PAGES	STATUS	STOP
STORAGE	STORED (R)	STRAIGHT_JOIN (R)
STRING	SUBCLASS_ORIGIN	SUBJECT
SUBPARTITION	SUBPARTITIONS	SUPER
SUSPEND	SWAPS	SWITCHES
TABLE (R)	TABLES	TABLESPACE
TABLE_CHECKSUM	TABLE_NAME	TEMPORARY
TEMPTABLE	TERMINATED (R)	TEXT
THAN	THEN (R)	TIES[ak]
TIME	TIMESTAMP	TIMESTAMPADD
TIMESTAMPDIFF	TINYBLOB (R)	TINYINT (R)
TINYTEXT (R)	TO (R)	TRAILING (R)
TRANSACTION	TRIGGER (R)	TRIGGERS
TRUE (R)	TRUNCATE	TYPE
TYPES	UNBOUNDED[al]	UNCOMMITTED
UNDEFINED	UNDO (R)	UNDOFILE
UNDO_BUFFER_SIZE	UNICODE	UNINSTALL
UNION (R)	UNIQUE (R)	UNKNOWN
UNLOCK (R)	UNSIGNED (R)	UNTIL
UPDATE (R)	UPGRADE	USAGE (R)

USE (R)	USER	USER_RESOURCES
USE_FRM	USING (R)	UTC_DATE (R)
UTC_TIME (R)	UTC_TIMESTAMP (R)	VALIDATION
VALUE	VALUES (R)	VARBINARY (R)
VARCHAR (R)	VARCHARACTER (R)	VARIABLES
VARYING (R)	VIEW	VIRTUAL (R)
VISIBLE	WAIT	WARNINGS
WEEK	WEIGHT_STRING	WHEN (R)
WHERE (R)	WHILE (R)	WINDOW[am] (R)
WITH (R)	WITHOUT	WORK
WRAPPER	WRITE (R)	X509
XA	XID	XML
XOR (R)	YEAR	YEAR_MONTH (R)
ZEROFILL (R)		

[a] ANALYSE: removed in 8.0.1
[b] BUCKETS: added in 8.0.2 (nonreserved)
[c] CLONE: added in 8.0.2 (nonreserved)
[d] CUBE: became reserved in 8.0.1
[e] CUME_DIST: added in 8.0.2 (reserved)
[f] DENSE_RANK: added in 8.0.2 (reserved)
[g] EXCLUDE: added in 8.0.2 (nonreserved)
[h] FIRST_VALUE: added in 8.0.2 (reserved)
[i] FOLLOWING: added in 8.0.2 (nonreserved)
[j] FUNCTION: became reserved in 8.0.1
[k] GROUPING: added in 8.0.1 (reserved)
[l] GROUPS: added in 8.0.2 (reserved)
[m] HISTOGRAM: added in 8.0.2 (nonreserved)
[n] LAG: added in 8.0.2 (reserved)
[o] LAST_VALUE: added in 8.0.2 (reserved)
[p] LEAD: added in 8.0.2 (reserved)
[q] LOCKED: added in 8.0.1 (nonreserved)
[r] NOWAIT: added in 8.0.1 (nonreserved)
[s] NTH_VALUE: added in 8.0.2 (reserved)
[t] NTILE: added in 8.0.2 (reserved)
[u] NULLS: added in 8.0.2 (nonreserved)
[v] OF: added in 8.0.1 (reserved)
[w] OTHERS: added in 8.0.2 (nonreserved)
[x] OVER: added in 8.0.2 (reserved)
[y] PERCENT_RANK: added in 8.0.2 (reserved)
[z] PERSIST_ONLY: added in 8.0.2 (reserved)
[aa] PRECEDING: added in 8.0.2 (nonreserved)
[ab] RANK: added in 8.0.2 (reserved)
[ac] RECURSIVE: added in 8.0.1 (reserved)
[ad] REMOTE: added in 8.0.2 (nonreserved)
[ae] RESPECT: added in 8.0.2 (nonreserved)
[af] ROLE: became nonreserved in 8.0.1
[ag] ROW: became reserved in 8.0.2
[ah] ROWS: became reserved in 8.0.2
[ai] ROW_NUMBER: added in 8.0.2 (reserved)
[aj] SKIP: added in 8.0.1 (nonreserved)
[ak] TIES: added in 8.0.2 (nonreserved)
[al] UNBOUNDED: added in 8.0.2 (nonreserved)
[am] WINDOW: added in 8.0.2 (reserved)

The following table shows the keywords and reserved words that are added in MySQL 8.0. Reserved keywords are marked with (R).

Table 9.3 Keywords and Reserved Words Added in MySQL 8.0 Compared to MySQL 5.7

ADMIN (R)	BUCKETS	CLONE
COMPONENT	CUME_DIST (R)	DENSE_RANK (R)
EXCEPT (R)	EXCLUDE	FIRST_VALUE (R)
FOLLOWING	GROUPING (R)	GROUPS (R)
HISTOGRAM	INVISIBLE	LAG (R)
LAST_VALUE (R)	LEAD (R)	LOCKED
NOWAIT	NTH_VALUE (R)	NTILE (R)
NULLS	OF (R)	OTHERS
OVER (R)	PERCENT_RANK (R)	PERSIST (R)
PERSIST_ONLY (R)	PRECEDING	RANK (R)
RECURSIVE (R)	REMOTE	RESPECT
ROLE	ROW_NUMBER (R)	SKIP
TIES	UNBOUNDED	VISIBLE
WINDOW (R)		

The following table shows the keywords and reserved words that are removed in MySQL 8.0. Reserved keywords are marked with (R).

Table 9.4 Keywords and Reserved Words Removed in MySQL 8.0 Compared to MySQL 5.7

ANALYSE	PARSE_GCOL_EXPR	

9.4 User-Defined Variables

You can store a value in a user-defined variable in one statement and refer to it later in another statement. This enables you to pass values from one statement to another.

User variables are written as *@var_name*, where the variable name *var_name* consists of alphanumeric characters, ., _, and $. A user variable name can contain other characters if you quote it as a string or identifier (for example, @'my-var', @"my-var", or @`my-var`).

User-defined variables are session specific. A user variable defined by one client cannot be seen or used by other clients. (Exception: A user with access to the Performance Schema user_variables_by_thread table can see all user variables for all sessions.) All variables for a given client session are automatically freed when that client exits.

User variable names are not case sensitive. Names have a maximum length of 64 characters.

One way to set a user-defined variable is by issuing a SET statement:

```
SET @var_name = expr [, @var_name = expr] ...
```

For SET, either = or := can be used as the assignment operator.

You can also assign a value to a user variable in statements other than SET. In this case, the assignment operator must be := and not = because the latter is treated as the comparison operator = in non-SET statements:

```
mysql> SET @t1=1, @t2=2, @t3:=4;
mysql> SELECT @t1, @t2, @t3, @t4 := @t1+@t2+@t3;
+------+------+------+--------------------+
```

```
| @t1  | @t2  | @t3  | @t4 := @t1+@t2+@t3 |
+------+------+------+--------------------+
|    1 |    2 |    4 |                  7 |
+------+------+------+--------------------+
```

User variables can be assigned a value from a limited set of data types: integer, decimal, floating-point, binary or nonbinary string, or `NULL` value. Assignment of decimal and real values does not preserve the precision or scale of the value. A value of a type other than one of the permissible types is converted to a permissible type. For example, a value having a temporal or spatial data type is converted to a binary string. A value having the `JSON` data type is converted to a string with a character set of `utf8mb4` and a collation of `utf8mb4_bin`.

If a user variable is assigned a nonbinary (character) string value, it has the same character set and collation as the string. The coercibility of user variables is implicit. (This is the same coercibility as for table column values.)

Hexadcimal or bit values assigned to user variables are treated as binary strings. To assign a hexadecimal or bit value as a number to a user variable, use it in numeric context. For example, add 0 or use `CAST(... AS UNSIGNED)`:

```
mysql> SET @v1 = X'41';
mysql> SET @v2 = X'41'+0;
mysql> SET @v3 = CAST(X'41' AS UNSIGNED);
mysql> SELECT @v1, @v2, @v3;
+------+------+------+
| @v1  | @v2  | @v3  |
+------+------+------+
| A    |   65 |   65 |
+------+------+------+
mysql> SET @v1 = b'1000001';
mysql> SET @v2 = b'1000001'+0;
mysql> SET @v3 = CAST(b'1000001' AS UNSIGNED);
mysql> SELECT @v1, @v2, @v3;
+------+------+------+
| @v1  | @v2  | @v3  |
+------+------+------+
| A    |   65 |   65 |
+------+------+------+
```

If the value of a user variable is selected in a result set, it is returned to the client as a string.

If you refer to a variable that has not been initialized, it has a value of `NULL` and a type of string.

User variables may be used in most contexts where expressions are permitted. This does not currently include contexts that explicitly require a literal value, such as in the `LIMIT` clause of a `SELECT` statement, or the `IGNORE N LINES` clause of a `LOAD DATA` statement.

As a general rule, other than in `SET` statements, you should never assign a value to a user variable and read the value within the same statement. For example, to increment a variable, this is okay:

```
SET @a = @a + 1;
```

For other statements, such as `SELECT`, you might get the results you expect, but this is not guaranteed. In the following statement, you might think that MySQL will evaluate `@a` first and then do an assignment second:

```
SELECT @a, @a:=@a+1, ...;
```

However, the order of evaluation for expressions involving user variables is undefined.

Another issue with assigning a value to a variable and reading the value within the same non-`SET` statement is that the default result type of a variable is based on its type at the start of the statement. The following example illustrates this:

```
mysql> SET @a='test';
mysql> SELECT @a,(@a:=20) FROM tbl_name;
```

For this SELECT statement, MySQL reports to the client that column one is a string and converts all accesses of @a to strings, even though @a is set to a number for the second row. After the SELECT statement executes, @a is regarded as a number for the next statement.

To avoid problems with this behavior, either do not assign a value to and read the value of the same variable within a single statement, or else set the variable to 0, 0.0, or '' to define its type before you use it.

In a SELECT statement, each select expression is evaluated only when sent to the client. This means that in a HAVING, GROUP BY, or ORDER BY clause, referring to a variable that is assigned a value in the select expression list does *not* work as expected:

```
mysql> SELECT (@aa:=id) AS a, (@aa+3) AS b FROM tbl_name HAVING b=5;
```

The reference to b in the HAVING clause refers to an alias for an expression in the select list that uses @aa. This does not work as expected: @aa contains the value of id from the previous selected row, not from the current row.

User variables are intended to provide data values. They cannot be used directly in an SQL statement as an identifier or as part of an identifier, such as in contexts where a table or database name is expected, or as a reserved word such as SELECT. This is true even if the variable is quoted, as shown in the following example:

```
mysql> SELECT c1 FROM t;
+----+
| c1 |
+----+
|  0 |
+----+
|  1 |
+----+
2 rows in set (0.00 sec)

mysql> SET @col = "c1";
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @col FROM t;
+------+
| @col |
+------+
| c1   |
+------+
1 row in set (0.00 sec)

mysql> SELECT `@col` FROM t;
ERROR 1054 (42S22): Unknown column '@col' in 'field list'

mysql> SET @col = "`c1`";
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @col FROM t;
+------+
| @col |
+------+
| `c1` |
+------+
1 row in set (0.00 sec)
```

An exception to this principle that user variables cannot be used to provide identifiers, is when you are constructing a string for use as a prepared statement to execute later. In this case, user variables can be used to provide any part of the statement. The following example illustrates how this can be done:

```
mysql> SET @c = "c1";
Query OK, 0 rows affected (0.00 sec)

mysql> SET @s = CONCAT("SELECT ", @c, " FROM t");
Query OK, 0 rows affected (0.00 sec)

mysql> PREPARE stmt FROM @s;
Query OK, 0 rows affected (0.04 sec)
Statement prepared

mysql> EXECUTE stmt;
+----+
| c1 |
+----+
|  0 |
+----+
|  1 |
+----+
2 rows in set (0.00 sec)

mysql> DEALLOCATE PREPARE stmt;
Query OK, 0 rows affected (0.00 sec)
```

See Section 13.5, “Prepared SQL Statement Syntax”, for more information.

A similar technique can be used in application programs to construct SQL statements using program variables, as shown here using PHP 5:

```
<?php
  $mysqli = new mysqli("localhost", "user", "pass", "test");

  if( mysqli_connect_errno() )
    die("Connection failed: %s\n", mysqli_connect_error());

  $col = "c1";

  $query = "SELECT $col FROM t";

  $result = $mysqli->query($query);

  while($row = $result->fetch_assoc())
  {
    echo "<p>" . $row["$col"] . "</p>\n";
  }

  $result->close();

  $mysqli->close();
?>
```

Assembling an SQL statement in this fashion is sometimes known as “Dynamic SQL”.

9.5 Expression Syntax

The following rules define expression syntax in MySQL. The grammar shown here is based on that given in the `sql/sql_yacc.yy` file of MySQL source distributions. See the notes after the grammar for additional information about some of the terms.

```
expr:
    expr OR expr
  | expr || expr
  | expr XOR expr
  | expr AND expr
  | expr && expr
  | NOT expr
  | ! expr
```

```
  | boolean_primary IS [NOT] {TRUE | FALSE | UNKNOWN}
  | boolean_primary

boolean_primary:
    boolean_primary IS [NOT] NULL
  | boolean_primary <=> predicate
  | boolean_primary comparison_operator predicate
  | boolean_primary comparison_operator {ALL | ANY} (subquery)
  | predicate

comparison_operator: = | >= | > | <= | < | <> | !=

predicate:
    bit_expr [NOT] IN (subquery)
  | bit_expr [NOT] IN (expr [, expr] ...)
  | bit_expr [NOT] BETWEEN bit_expr AND predicate
  | bit_expr SOUNDS LIKE bit_expr
  | bit_expr [NOT] LIKE simple_expr [ESCAPE simple_expr]
  | bit_expr [NOT] REGEXP bit_expr
  | bit_expr

bit_expr:
    bit_expr | bit_expr
  | bit_expr & bit_expr
  | bit_expr << bit_expr
  | bit_expr >> bit_expr
  | bit_expr + bit_expr
  | bit_expr - bit_expr
  | bit_expr * bit_expr
  | bit_expr / bit_expr
  | bit_expr DIV bit_expr
  | bit_expr MOD bit_expr
  | bit_expr % bit_expr
  | bit_expr ^ bit_expr
  | bit_expr + interval_expr
  | bit_expr - interval_expr
  | simple_expr

simple_expr:
    literal
  | identifier
  | function_call
  | simple_expr COLLATE collation_name
  | param_marker
  | variable
  | simple_expr || simple_expr
  | + simple_expr
  | - simple_expr
  | ~ simple_expr
  | ! simple_expr
  | BINARY simple_expr
  | (expr [, expr] ...)
  | ROW (expr, expr [, expr] ...)
  | (subquery)
  | EXISTS (subquery)
  | {identifier expr}
  | match_expr
  | case_expr
  | interval_expr
```

Notes:

For operator precedence, see in Section 12.3.1, “Operator Precedence”.

For literal value syntax, see Section 9.1, “Literal Values”.

For identifier syntax, see Section 9.2, “Schema Object Names”.

Variables can be user variables, system variables, or stored program local variables or parameters:

- User variables: Section 9.4, “User-Defined Variables”

- System variables: Section 5.1.6, “Using System Variables”
- Local variables: Section 13.6.4.1, “Local Variable DECLARE Syntax”
- Parameters: Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”

param_marker is ? as used in prepared statements for placeholders. See Section 13.5.1, “PREPARE Syntax”.

(*subquery*) indicates a subquery that returns a single value; that is, a scalar subquery. See Section 13.2.11.1, “The Subquery as Scalar Operand”.

{*identifier expr*} is ODBC escape syntax and is accepted for ODBC compatibility. The value is *expr*. The curly braces in the syntax should be written literally; they are not metasyntax as used elsewhere in syntax descriptions.

match_expr indicates a MATCH expression. See Section 12.9, “Full-Text Search Functions”.

case_expr indicates a CASE expression. See Section 12.4, “Control Flow Functions”.

interval_expr represents a time interval. The syntax is INTERVAL *expr unit*, where *unit* is a specifier such as HOUR, DAY, or WEEK. For the full list of *unit* specifiers, see the description of the DATE_ADD() function in Section 12.7, “Date and Time Functions”.

The meaning of some operators depends on the SQL mode:

- By default, || is a logical OR operator. With PIPES_AS_CONCAT enabled, || is string concatenation, with a precedence between ^ and the unary operators.
- By default, ! has a higher precedence than NOT. With HIGH_NOT_PRECEDENCE enabled, ! and NOT have the same precedence.

See Section 5.1.8, “Server SQL Modes”.

9.6 Comment Syntax

MySQL Server supports three comment styles:

- From a # character to the end of the line.
- From a -- sequence to the end of the line. In MySQL, the -- (double-dash) comment style requires the second dash to be followed by at least one whitespace or control character (such as a space, tab, newline, and so on). This syntax differs slightly from standard SQL comment syntax, as discussed in Section 1.8.2.4, “'--' as the Start of a Comment”.
- From a /* sequence to the following */ sequence, as in the C programming language. This syntax enables a comment to extend over multiple lines because the beginning and closing sequences need not be on the same line.

The following example demonstrates all three comment styles:

```
mysql> SELECT 1+1;     # This comment continues to the end of line
mysql> SELECT 1+1;     -- This comment continues to the end of line
mysql> SELECT 1 /* this is an in-line comment */ + 1;
mysql> SELECT 1+
/*
this is a
multiple-line comment
*/
1;
```

Nested comments are not supported. (Under some conditions, nested comments might be permitted, but usually are not, and users should avoid them.)

MySQL Server supports some variants of C-style comments. These enable you to write code that includes MySQL extensions, but is still portable, by using comments of the following form:

```
/*! MySQL-specific code */
```

In this case, MySQL Server parses and executes the code within the comment as it would any other SQL statement, but other SQL servers will ignore the extensions. For example, MySQL Server recognizes the `STRAIGHT_JOIN` keyword in the following statement, but other servers will not:

```
SELECT /*! STRAIGHT_JOIN */ col1 FROM table1,table2 WHERE ...
```

If you add a version number after the `!` character, the syntax within the comment is executed only if the MySQL version is greater than or equal to the specified version number. The `KEY_BLOCK_SIZE` keyword in the following comment is executed only by servers from MySQL 5.1.10 or higher:

```
CREATE TABLE t1(a INT, KEY (a)) /*!50110 KEY_BLOCK_SIZE=1024 */;
```

The comment syntax just described applies to how the `mysqld` server parses SQL statements. The `mysql` client program also performs some parsing of statements before sending them to the server. (It does this to determine statement boundaries within a multiple-statement input line.)

Comments in this format, `/*!12345 ... */`, are not stored on the server. If this format is used to comment stored routines, the comments will not be retained on the server.

Another variant of C-style comment syntax is used to specify optimizer hints. Hint comments include a `+` character following the `/*` comment opening sequence. Example:

```
SELECT /*+ BKA(t1) */ FROM ... ;
```

For more information, see Section 8.9.2, “Optimizer Hints”.

The use of short-form `mysql` commands such as `\C` within multiple-line `/* ... */` comments is not supported.

Chapter 10 Globalization

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This chapter covers issues of globalization, which includes internationalization (MySQL's capabilities for adapting to local use) and localization (selecting particular local conventions):

- MySQL support for character sets in SQL statements.
- How to configure the server to support different character sets.
- Selecting the language for error messages.
- How to set the server's time zone and enable per-connection time zone support.
- Selecting the locale for day and month names.

10.1 Character Set Support

MySQL includes character set support that enables you to store data using a variety of character sets and perform comparisons according to a variety of collations. You can specify character sets at the server, database, table, and column level. MySQL supports the use of character sets for the `MyISAM`, `MEMORY`, and `InnoDB` storage engines.

This chapter discusses the following topics:

- What are character sets and collations?

- The multiple-level default system for character set assignment.
- Syntax for specifying character sets and collations.
- Affected functions and operations.
- Unicode support.
- The character sets and collations that are available, with notes.

Character set issues affect not only data storage, but also communication between client programs and the MySQL server. If you want the client program to communicate with the server using a character set different from the default, you'll need to indicate which one. For example, to use the `utf8` Unicode character set, issue this statement after connecting to the server:

```
SET NAMES 'utf8';
```

For more information about configuring character sets for application use and character set-related issues in client/server communication, see Section 10.1.5, “Configuring Application Character Set and Collation”, and Section 10.1.4, “Connection Character Sets and Collations”.

10.1.1 Character Sets and Collations in General

A *character set* is a set of symbols and encodings. A *collation* is a set of rules for comparing characters in a character set. Let's make the distinction clear with an example of an imaginary character set.

Suppose that we have an alphabet with four letters: A, B, a, b. We give each letter a number: A = 0, B = 1, a = 2, b = 3. The letter A is a symbol, the number 0 is the *encoding* for A, and the combination of all four letters and their encodings is a *character set*.

Suppose that we want to compare two string values, A and B. The simplest way to do this is to look at the encodings: 0 for A and 1 for B. Because 0 is less than 1, we say A is less than B. What we've just done is apply a collation to our character set. The collation is a set of rules (only one rule in this case): “compare the encodings.” We call this simplest of all possible collations a *binary* collation.

But what if we want to say that the lowercase and uppercase letters are equivalent? Then we would have at least two rules: (1) treat the lowercase letters a and b as equivalent to A and B; (2) then compare the encodings. We call this a *case-insensitive* collation. It is a little more complex than a binary collation.

In real life, most character sets have many characters: not just A and B but whole alphabets, sometimes multiple alphabets or eastern writing systems with thousands of characters, along with many special symbols and punctuation marks. Also in real life, most collations have many rules, not just for whether to distinguish lettercase, but also for whether to distinguish accents (an “accent” is a mark attached to a character as in German Ö), and for multiple-character mappings (such as the rule that Ö = OE in one of the two German collations).

MySQL can do these things for you:

- Store strings using a variety of character sets.
- Compare strings using a variety of collations.
- Mix strings with different character sets or collations in the same server, the same database, or even the same table.
- Enable specification of character set and collation at any level.

To use these features effectively, you must know what character sets and collations are available, how to change the defaults, and how they affect the behavior of string operators and functions.

10.1.2 Character Sets and Collations in MySQL

MySQL Server supports multiple character sets. To display the available character sets, use the `INFORMATION_SCHEMA CHARACTER_SETS` table or the `SHOW CHARACTER SET` statement. A partial listing follows. For more complete information, see Section 10.1.10, “Supported Character Sets and Collations”.

```
mysql> SHOW CHARACTER SET;
+----------+---------------------------------+---------------------+--------+
| Charset  | Description                     | Default collation   | Maxlen |
+----------+---------------------------------+---------------------+--------+
| big5     | Big5 Traditional Chinese        | big5_chinese_ci     |      2 |
...
| latin1   | cp1252 West European            | latin1_swedish_ci   |      1 |
| latin2   | ISO 8859-2 Central European     | latin2_general_ci   |      1 |
...
| utf8     | UTF-8 Unicode                   | utf8_general_ci     |      3 |
| ucs2     | UCS-2 Unicode                   | ucs2_general_ci     |      2 |
...
| utf8mb4  | UTF-8 Unicode                   | utf8mb4_general_ci  |      4 |
...
| binary   | Binary pseudo charset           | binary              |      1 |
...
```

By default, the `SHOW CHARACTER SET` statement displays all available character sets. It takes an optional `LIKE` or `WHERE` clause that indicates which character set names to match. For example:

```
mysql> SHOW CHARACTER SET LIKE 'latin%';
+---------+----------------------------+-------------------+--------+
| Charset | Description                | Default collation | Maxlen |
+---------+----------------------------+-------------------+--------+
| latin1  | cp1252 West European       | latin1_swedish_ci |      1 |
| latin2  | ISO 8859-2 Central European | latin2_general_ci |      1 |
| latin5  | ISO 8859-9 Turkish         | latin5_turkish_ci |      1 |
| latin7  | ISO 8859-13 Baltic         | latin7_general_ci |      1 |
+---------+----------------------------+-------------------+--------+
```

A given character set always has at least one collation, and most character sets have several. To list the display collations for a character set, use the `INFORMATION_SCHEMA COLLATIONS` table or the `SHOW COLLATION` statement.

By default, the `SHOW COLLATION` statement displays all available collations. It takes an optional `LIKE` or `WHERE` clause that indicates which collation names to display. For example, to see the collations for the `latin1` (cp1252 West European) character set, use this statement:

```
mysql> SHOW COLLATION WHERE Charset = 'latin1';
+-------------------+---------+----+---------+----------+---------+
| Collation         | Charset | Id | Default | Compiled | Sortlen |
+-------------------+---------+----+---------+----------+---------+
| latin1_german1_ci | latin1  |  5 |         | Yes      |       1 |
| latin1_swedish_ci | latin1  |  8 | Yes     | Yes      |       1 |
| latin1_danish_ci  | latin1  | 15 |         | Yes      |       1 |
| latin1_german2_ci | latin1  | 31 |         | Yes      |       2 |
| latin1_bin        | latin1  | 47 |         | Yes      |       1 |
| latin1_general_ci | latin1  | 48 |         | Yes      |       1 |
| latin1_general_cs | latin1  | 49 |         | Yes      |       1 |
| latin1_spanish_ci | latin1  | 94 |         | Yes      |       1 |
+-------------------+---------+----+---------+----------+---------+
```

The `latin1` collations have the following meanings.

Collation	Meaning
`latin1_bin`	Binary according to `latin1` encoding
`latin1_danish_ci`	Danish/Norwegian

Collation	Meaning
latin1_general_ci	Multilingual (Western European)
latin1_general_cs	Multilingual (ISO Western European), case sensitive
latin1_german1_ci	German DIN-1 (dictionary order)
latin1_german2_ci	German DIN-2 (phone book order)
latin1_spanish_ci	Modern Spanish
latin1_swedish_ci	Swedish/Finnish

Collations have these general characteristics:

- Two different character sets cannot have the same collation.
- Each character set has a *default collation*. For example, the default collations for latin1 and utf8 are latin1_swedish_ci and utf8_general_ci, respectively. The INFORMATION_SCHEMA CHARACTER_SETS table and the SHOW CHARACTER SET statement indicate the default collation for each character set. The INFORMATION_SCHEMA COLLATIONS table and the SHOW COLLATION statement have a column that indicates for each collation whether it is the default for its character set (Yes if so, empty if not).
- Collation names start with the name of the character set with which they are associated, generally followed by one or more suffixes indicating other collation characteristics. For additional information about naming conventions, see Section 10.1.3.1, “Collation Naming Conventions”.

When a character set has multiple collations, it might not be clear which collation is most suitable for a given application. To avoid choosing an inappropriate collation, perform some comparisons with representative data values to make sure that a given collation sorts values the way you expect.

10.1.2.1 Character Set Repertoire

The *repertoire* of a character set is the collection of characters in the set.

String expressions have a repertoire attribute, which can have two values:

- ASCII: The expression can contain only characters in the Unicode range U+0000 to U+007F.
- UNICODE: The expression can contain characters in the Unicode range U+0000 to U+10FFFF. This includes characters in the Basic Multilingual Plane (BMP) range (U+0000 to U+FFFF) and supplementary characters outside the BMP range (U+01000 to U+10FFFF).

The ASCII range is a subset of UNICODE range, so a string with ASCII repertoire can be converted safely without loss of information to the character set of any string with UNICODE repertoire or to a character set that is a superset of ASCII. (All MySQL character sets are supersets of ASCII with the exception of swe7, which reuses some punctuation characters for Swedish accented characters.) The use of repertoire enables character set conversion in expressions for many cases where MySQL would otherwise return an “illegal mix of collations” error.

The following discussion provides examples of expressions and their repertoires, and describes how the use of repertoire changes string expression evaluation:

- The repertoire for a string constant depends on string content and may differ from the repertoire of the string character set. Consider these statements:

```
SET NAMES utf8; SELECT 'abc';
SELECT _utf8'def';
SELECT N'MySQL';
```

  Although the character set is utf8 in each of the preceding cases, the strings do not actually contain any characters outside the ASCII range, so their repertoire is ASCII rather than UNICODE.

- A column having the `ascii` character set has `ASCII` repertoire because of its character set. In the following table, `c1` has `ASCII` repertoire:

```
CREATE TABLE t1 (c1 CHAR(1) CHARACTER SET ascii);
```

The following example illustrates how repertoire enables a result to be determined in a case where an error occurs without repertoire:

```
CREATE TABLE t1 (
  c1 CHAR(1) CHARACTER SET latin1,
  c2 CHAR(1) CHARACTER SET ascii
);
INSERT INTO t1 VALUES ('a','b');
SELECT CONCAT(c1,c2) FROM t1;
```

Without repertoire, this error occurs:

```
ERROR 1267 (HY000): Illegal mix of collations (latin1_swedish_ci,IMPLICIT)
and (ascii_general_ci,IMPLICIT) for operation 'concat'
```

Using repertoire, subset to superset (`ascii` to `latin1`) conversion can occur and a result is returned:

```
+---------------+
| CONCAT(c1,c2) |
+---------------+
| ab            |
+---------------+
```

- Functions with one string argument inherit the repertoire of their argument. The result of `UPPER(_utf8'abc')` has `ASCII` repertoire because its argument has `ASCII` repertoire.
- For functions that return a string but do not have string arguments and use `character_set_connection` as the result character set, the result repertoire is `ASCII` if `character_set_connection` is `ascii`, and `UNICODE` otherwise:

```
FORMAT(numeric_column, 4);
```

Use of repertoire changes how MySQL evaluates the following example:

```
SET NAMES ascii;
CREATE TABLE t1 (a INT, b VARCHAR(10) CHARACTER SET latin1);
INSERT INTO t1 VALUES (1,'b');
SELECT CONCAT(FORMAT(a, 4), b) FROM t1;
```

Without repertoire, this error occurs:

```
ERROR 1267 (HY000): Illegal mix of collations (ascii_general_ci,COERCIBLE)
and (latin1_swedish_ci,IMPLICIT) for operation 'concat'
```

With repertoire, a result is returned:

```
+-------------------------+
| CONCAT(FORMAT(a, 4), b) |
+-------------------------+
| 1.0000b                 |
+-------------------------+
```

- Functions with two or more string arguments use the “widest” argument repertoire for the result repertoire (`UNICODE` is wider than `ASCII`). Consider the following `CONCAT()` calls:

```
CONCAT(_ucs2 X'0041', _ucs2 X'0042')
CONCAT(_ucs2 X'0041', _ucs2 X'00C2')
```

For the first call, the repertoire is `ASCII` because both arguments are within the range of the `ascii` character set. For the second call, the repertoire is `UNICODE` because the second argument is outside the `ascii` character set range.

- The repertoire for function return values is determined based only on the repertoire of the arguments that affect the result's character set and collation.

```
IF(column1 < column2, 'smaller', 'greater')
```

The result repertoire is `ASCII` because the two string arguments (the second argument and the third argument) both have `ASCII` repertoire. The first argument does not matter for the result repertoire, even if the expression uses string values.

10.1.2.2 UTF-8 for Metadata

Metadata is “the data about the data.” Anything that *describes* the database—as opposed to being the *contents* of the database—is metadata. Thus column names, database names, user names, version names, and most of the string results from `SHOW` are metadata. This is also true of the contents of tables in `INFORMATION_SCHEMA` because those tables by definition contain information about database objects.

Representation of metadata must satisfy these requirements:

- All metadata must be in the same character set. Otherwise, neither the `SHOW` statements nor `SELECT` statements for tables in `INFORMATION_SCHEMA` would work properly because different rows in the same column of the results of these operations would be in different character sets.
- Metadata must include all characters in all languages. Otherwise, users would not be able to name columns and tables using their own languages.

To satisfy both requirements, MySQL stores metadata in a Unicode character set, namely UTF-8. This does not cause any disruption if you never use accented or non-Latin characters. But if you do, you should be aware that metadata is in UTF-8.

The metadata requirements mean that the return values of the `USER()`, `CURRENT_USER()`, `SESSION_USER()`, `SYSTEM_USER()`, `DATABASE()`, and `VERSION()` functions have the UTF-8 character set by default.

The server sets the `character_set_system` system variable to the name of the metadata character set:

```
mysql> SHOW VARIABLES LIKE 'character_set_system';
+----------------------+-------+
| Variable_name        | Value |
+----------------------+-------+
| character_set_system | utf8  |
+----------------------+-------+
```

Storage of metadata using Unicode does *not* mean that the server returns headers of columns and the results of `DESCRIBE` functions in the `character_set_system` character set by default. When you use `SELECT column1 FROM t`, the name `column1` itself is returned from the server to the client in the character set determined by the value of the `character_set_results` system variable, which has a default value of `latin1`. If you want the server to pass metadata results back in a different character set, use the `SET NAMES` statement to force the server to perform character set conversion. `SET NAMES` sets the `character_set_results` and other related system variables. (See Section 10.1.4, “Connection Character Sets and Collations”.) Alternatively, a client program can

perform the conversion after receiving the result from the server. It is more efficient for the client to perform the conversion, but this option is not always available for all clients.

If `character_set_results` is set to `NULL`, no conversion is performed and the server returns metadata using its original character set (the set indicated by `character_set_system`).

Error messages returned from the server to the client are converted to the client character set automatically, as with metadata.

If you are using (for example) the `USER()` function for comparison or assignment within a single statement, don't worry. MySQL performs some automatic conversion for you.

```
SELECT * FROM t1 WHERE USER() = latin1_column;
```

This works because the contents of `latin1_column` are automatically converted to UTF-8 before the comparison.

```
INSERT INTO t1 (latin1_column) SELECT USER();
```

This works because the contents of `USER()` are automatically converted to `latin1` before the assignment.

Although automatic conversion is not in the SQL standard, the standard does say that every character set is (in terms of supported characters) a “subset” of Unicode. Because it is a well-known principle that “what applies to a superset can apply to a subset,” we believe that a collation for Unicode can apply for comparisons with non-Unicode strings. For more information about coercion of strings, see Section 10.1.8.4, “Collation Coercibility in Expressions”.

10.1.3 Specifying Character Sets and Collations

There are default settings for character sets and collations at four levels: server, database, table, and column. The description in the following sections may appear complex, but it has been found in practice that multiple-level defaulting leads to natural and obvious results.

`CHARACTER SET` is used in clauses that specify a character set. `CHARSET` can be used as a synonym for `CHARACTER SET`.

Character set issues affect not only data storage, but also communication between client programs and the MySQL server. If you want the client program to communicate with the server using a character set different from the default, you'll need to indicate which one. For example, to use the `utf8mb4` Unicode character set, issue this statement after connecting to the server:

```
SET NAMES 'utf8mb4';
```

For more information about character set-related issues in client/server communication, see Section 10.1.4, “Connection Character Sets and Collations”.

10.1.3.1 Collation Naming Conventions

MySQL collation names follow these conventions:

- A collation name starts with the name of the character set with which it is associated, generally followed by one or more suffixes indicating other collation characteristics. For example, `utf8mb4_general_ci` and `latin1_swedish_ci` are collations for the `utf8mb4` and `latin1` character sets, respectively. The `binary` character set has a single collation, also named `binary`, with no suffixes.

- A language-specific collation includes a locale code or language name. For example, `utf8mb4_tr_0900_ai_ci` and `utf8mb4_hu_0900_ai_ci` sort characters for the `utf8mb4`

character set using the rules of Turkish and Hungarian, respectively. `utf8mb4_turkish_ci` and `utf8mb4_hungarian_ci` are similar but based on a less recent version of the Unicode Collation Algorithm.

- Collation suffixes indicate whether a collation is case and accent sensitive, or binary. The following table shows the suffixes used to indicate these characteristics.

  Table 10.1 Collation Case/Accent Sensitivity Suffixes

Suffix	Meaning
`_ai`	Accent insensitive
`_as`	Accent sensitive
`_ci`	Case insensitive
`_cs`	Case sensitive
`_ks`	Kana sensitive
`_bin`	Binary

  For nonbinary collation names that do not specify accent sensitivity, it is determined by case sensitivity. If a collation name does not contain `_ai` or `_as`, `_ci` in the name implies `_ai` and `_cs` in the name implies `_as`. For example, `latin1_general_ci` is explicitly case insensitive and implicitly accent insensitive, `latin1_general_cs` is explicitly case sensitive and implicitly accent sensitive, and `utf8mb4_0900_ai_ci` is explicitly case and accent insensitive.

  For Japanese collations, the `_ks` suffix indicates that a collation is kana sensitive; that is, it distinguishes Katakana characters from Hiragana characters. Japanese collations without the `_ks` suffix are not kana sensitive and treat Katakana and Hiragana characters equal for sorting.

  For the `binary` collation of the `binary` character set, comparisons are based on numeric byte values. For the `_bin` collation of a nonbinary character set, comparisons are based on numeric character code values, which differ from byte values for multibyte characters. For more information, see Section 10.1.8.5, “The binary Collation Compared to _bin Collations”.

- For Unicode character sets, collation names may include a version number to indicate the version of the Unicode Collation Algorithm (UCA) on which the collation is based. UCA-based collations without a version number in the name use the version-4.0.0 UCA weight keys. For example:

  - `utf8mb4_0900_ai_ci` is based on UCA 9.0.0 weight keys (http://www.unicode.org/Public/UCA/9.0.0/allkeys.txt).
  - `utf8mb4_unicode_520_ci` is based on UCA 5.2.0 weight keys (http://www.unicode.org/Public/UCA/5.2.0/allkeys.txt).
  - `utf8mb4_unicode_ci` (with no version named) is based on UCA 4.0.0 weight keys (http://www.unicode.org/Public/UCA/4.0.0/allkeys-4.0.0.txt).

- For Unicode character sets, the *`xxx`*`_general_mysql500_ci` collations preserve the pre-5.1.24 ordering of the original *`xxx`*`_general_ci` collations and permit upgrades for tables created before MySQL 5.1.24 (Bug #27877).

10.1.3.2 Server Character Set and Collation

MySQL Server has a server character set and a server collation. These can be set at server startup on the command line or in an option file and changed at runtime.

Initially, the server character set and collation depend on the options that you use when you start `mysqld`. You can use `--character-set-server` for the character set. Along with it, you can add `--collation-server` for the collation. If you don't specify a character set, that is the same as saying `--character-set-server=latin1`. If you specify only a character set (for example, `latin1`) but

not a collation, that is the same as saying `--character-set-server=latin1 --collation-server=latin1_swedish_ci` because `latin1_swedish_ci` is the default collation for `latin1`. Therefore, the following three commands all have the same effect:

```
shell> mysqld
shell> mysqld --character-set-server=latin1
shell> mysqld --character-set-server=latin1 \
          --collation-server=latin1_swedish_ci
```

One way to change the settings is by recompiling. To change the default server character set and collation when building from sources, use the `DEFAULT_CHARSET` and `DEFAULT_COLLATION` options for `CMake`. For example:

```
shell> cmake . -DDEFAULT_CHARSET=latin1
```

Or:

```
shell> cmake . -DDEFAULT_CHARSET=latin1 \
          -DDEFAULT_COLLATION=latin1_german1_ci
```

Both `mysqld` and `CMake` verify that the character set/collation combination is valid. If not, each program displays an error message and terminates.

The server character set and collation are used as default values if the database character set and collation are not specified in `CREATE DATABASE` statements. They have no other purpose.

The current server character set and collation can be determined from the values of the `character_set_server` and `collation_server` system variables. These variables can be changed at runtime.

10.1.3.3 Database Character Set and Collation

Every database has a database character set and a database collation. The `CREATE DATABASE` and `ALTER DATABASE` statements have optional clauses for specifying the database character set and collation:

```
CREATE DATABASE db_name
    [[DEFAULT] CHARACTER SET charset_name]
    [[DEFAULT] COLLATE collation_name]

ALTER DATABASE db_name
    [[DEFAULT] CHARACTER SET charset_name]
    [[DEFAULT] COLLATE collation_name]
```

The keyword `SCHEMA` can be used instead of `DATABASE`.

The `CHARACTER SET` and `COLLATE` clauses make it possible to create databases with different character sets and collations on the same MySQL server.

Database options are stored in the data dictionary and can be examined by checking the `INFORMATION_SCHEMA.SCHEMATA` table.

Example:

```
CREATE DATABASE db_name CHARACTER SET latin1 COLLATE latin1_swedish_ci;
```

MySQL chooses the database character set and database collation in the following manner:

- If both `CHARACTER SET charset_name` and `COLLATE collation_name` are specified, character set `charset_name` and collation `collation_name` are used.

- If `CHARACTER SET charset_name` is specified without `COLLATE`, character set `charset_name` and its default collation are used. To see the default collation for each character set, use the `SHOW CHARACTER SET` statement.
- If `COLLATE collation_name` is specified without `CHARACTER SET`, the character set associated with `collation_name` and collation `collation_name` are used.
- Otherwise (neither `CHARACTER SET` nor `COLLATE` is specified), the server character set and server collation are used.

The character set and collation for the default database can be determined from the values of the `character_set_database` and `collation_database` system variables. The server sets these variables whenever the default database changes. If there is no default database, the variables have the same value as the corresponding server-level system variables, `character_set_server` and `collation_server`.

To see the default character set and collation for a given database, use these statements:

```
USE db_name;
SELECT @@character_set_database, @@collation_database;
```

Alternatively, to display the values without changing the default database:

```
SELECT DEFAULT_CHARACTER_SET_NAME, DEFAULT_COLLATION_NAME
FROM INFORMATION_SCHEMA.SCHEMATA WHERE SCHEMA_NAME = 'db_name';
```

The database character set and collation affect these aspects of server operation:

- For `CREATE TABLE` statements, the database character set and collation are used as default values for table definitions if the table character set and collation are not specified. To override this, provide explicit `CHARACTER SET` and `COLLATE` table options.
- For `LOAD DATA` statements that include no `CHARACTER SET` clause, the server uses the character set indicated by the `character_set_database` system variable to interpret the information in the file. To override this, provide an explicit `CHARACTER SET` clause.
- For stored routines (procedures and functions), the database character set and collation in effect at routine creation time are used as the character set and collation of character data parameters for which the declaration includes no `CHARACTER SET` or `COLLATE` attribute. To override this, provide explicit `CHARACTER SET` and `COLLATE` attributes.

10.1.3.4 Table Character Set and Collation

Every table has a table character set and a table collation. The `CREATE TABLE` and `ALTER TABLE` statements have optional clauses for specifying the table character set and collation:

```
CREATE TABLE tbl_name (column_list)
    [[DEFAULT] CHARACTER SET charset_name]
    [COLLATE collation_name]]

ALTER TABLE tbl_name
    [[DEFAULT] CHARACTER SET charset_name]
    [COLLATE collation_name]
```

Example:

```
CREATE TABLE t1 ( ... )
CHARACTER SET latin1 COLLATE latin1_danish_ci;
```

MySQL chooses the table character set and collation in the following manner:

- If both `CHARACTER SET charset_name` and `COLLATE collation_name` are specified, character set `charset_name` and collation `collation_name` are used.
- If `CHARACTER SET charset_name` is specified without `COLLATE`, character set `charset_name` and its default collation are used. To see the default collation for each character set, use the `SHOW CHARACTER SET` statement.
- If `COLLATE collation_name` is specified without `CHARACTER SET`, the character set associated with `collation_name` and collation `collation_name` are used.
- Otherwise (neither `CHARACTER SET` nor `COLLATE` is specified), the database character set and collation are used.

The table character set and collation are used as default values for column definitions if the column character set and collation are not specified in individual column definitions. The table character set and collation are MySQL extensions; there are no such things in standard SQL.

10.1.3.5 Column Character Set and Collation

Every “character” column (that is, a column of type `CHAR`, `VARCHAR`, or `TEXT`) has a column character set and a column collation. Column definition syntax for `CREATE TABLE` and `ALTER TABLE` has optional clauses for specifying the column character set and collation:

```
col_name {CHAR | VARCHAR | TEXT} (col_length)
    [CHARACTER SET charset_name]
    [COLLATE collation_name]
```

These clauses can also be used for `ENUM` and `SET` columns:

```
col_name {ENUM | SET} (val_list)
    [CHARACTER SET charset_name]
    [COLLATE collation_name]
```

Examples:

```
CREATE TABLE t1
(
    col1 VARCHAR(5)
      CHARACTER SET latin1
      COLLATE latin1_german1_ci
);

ALTER TABLE t1 MODIFY
    col1 VARCHAR(5)
      CHARACTER SET latin1
      COLLATE latin1_swedish_ci;
```

MySQL chooses the column character set and collation in the following manner:

- If both `CHARACTER SET charset_name` and `COLLATE collation_name` are specified, character set `charset_name` and collation `collation_name` are used.

  ```
  CREATE TABLE t1
  (
      col1 CHAR(10) CHARACTER SET utf8 COLLATE utf8_unicode_ci
  ) CHARACTER SET latin1 COLLATE latin1_bin;
  ```

  The character set and collation are specified for the column, so they are used. The column has character set `utf8` and collation `utf8_unicode_ci`.

- If `CHARACTER SET charset_name` is specified without `COLLATE`, character set `charset_name` and its default collation are used.

```
CREATE TABLE t1
(
    col1 CHAR(10) CHARACTER SET utf8
) CHARACTER SET latin1 COLLATE latin1_bin;
```

The character set is specified for the column, but the collation is not. The column has character set `utf8` and the default collation for `utf8`, which is `utf8_general_ci`. To see the default collation for each character set, use the `SHOW CHARACTER SET` statement.

- If `COLLATE collation_name` is specified without `CHARACTER SET`, the character set associated with `collation_name` and collation `collation_name` are used.

```
CREATE TABLE t1
(
    col1 CHAR(10) COLLATE utf8_polish_ci
) CHARACTER SET latin1 COLLATE latin1_bin;
```

The collation is specified for the column, but the character set is not. The column has collation `utf8_polish_ci` and the character set is the one associated with the collation, which is `utf8`.

- Otherwise (neither `CHARACTER SET` nor `COLLATE` is specified), the table character set and collation are used.

```
CREATE TABLE t1
(
    col1 CHAR(10)
) CHARACTER SET latin1 COLLATE latin1_bin;
```

Neither the character set nor collation is specified for the column, so the table defaults are used. The column has character set `latin1` and collation `latin1_bin`.

The `CHARACTER SET` and `COLLATE` clauses are standard SQL.

If you use `ALTER TABLE` to convert a column from one character set to another, MySQL attempts to map the data values, but if the character sets are incompatible, there may be data loss.

10.1.3.6 Character String Literal Character Set and Collation

Every character string literal has a character set and a collation.

For the simple statement `SELECT 'string'`, the string has the connection default character set and collation defined by the `character_set_connection` and `collation_connection` system variables.

A character string literal may have an optional character set introducer and `COLLATE` clause, to designate it as a string that uses a particular character set and collation:

```
[_charset_name]'string' [COLLATE collation_name]
```

Character set introducers and the `COLLATE` clause are implemented according to standard SQL specifications.

Examples:

```
SELECT 'abc';
SELECT _latin1'abc';
SELECT _binary'abc';
SELECT _utf8'abc' COLLATE utf8_danish_ci;
```

The `_charset_name` expression is formally called an *introducer*. It tells the parser, “the string that follows uses character set `charset_name`.” An introducer does not change the string to the introducer

character set like `CONVERT()` would do. It does not change the string value, although padding may occur. The introducer is just a signal.

MySQL determines the character set and collation of a character string literal in the following manner:

- If both *`_charset_name`* and `COLLATE` *`collation_name`* are specified, character set *`charset_name`* and collation *`collation_name`* are used. *`collation_name`* must be a permitted collation for *`charset_name`*.
- If *`_charset_name`* is specified but `COLLATE` is not specified, character set *`charset_name`* and its default collation are used. To see the default collation for each character set, use the `SHOW CHARACTER SET` statement.
- If *`_charset_name`* is not specified but `COLLATE` *`collation_name`* is specified, the connection default character set given by the `character_set_connection` system variable and collation *`collation_name`* are used. *`collation_name`* must be a permitted collation for the connection default character set.
- Otherwise (neither *`_charset_name`* nor `COLLATE` *`collation_name`* is specified), the connection default character set and collation given by the `character_set_connection` and `collation_connection` system variables are used.

Examples:

- A nonbinary string with `latin1` character set and `latin1_german1_ci` collation:

```
SELECT _latin1'Müller' COLLATE latin1_german1_ci;
```

- A nonbinary string with `utf8` character set and its default collation (that is, `utf8_general_ci`):

```
SELECT _utf8'Müller';
```

- A binary string with `binary` character set and its default collation (that is, `binary`):

```
SELECT _binary'Müller';
```

- A nonbinary string with the connection default character set and `utf8_general_ci` collation (fails if the connection character set is not `utf8`):

```
SELECT 'Müller' COLLATE utf8_general_ci;
```

- A string with the connection default character set and collation:

```
SELECT 'Müller';
```

An introducer indicates the character set for the following string, but does not change how the parser performs escape processing within the string. Escapes are always interpreted by the parser according to the character set given by `character_set_connection`.

The following examples show that escape processing occurs using `character_set_connection` even in the presence of an introducer. The examples use `SET NAMES` (which changes `character_set_connection`, as discussed in Section 10.1.4, “Connection Character Sets and Collations”), and display the resulting strings using the `HEX()` function so that the exact string contents can be seen.

Example 1:

```
mysql> SET NAMES latin1;
```

```
mysql> SELECT HEX('à\n'), HEX(_sjis'à\n');
+------------+-----------------+
| HEX('à\n')  | HEX(_sjis'à\n')  |
+------------+-----------------+
| E00A       | E00A            |
+------------+-----------------+
```

Here, `à` (hexadecimal value `E0`) is followed by `\n`, the escape sequence for newline. The escape sequence is interpreted using the `character_set_connection` value of `latin1` to produce a literal newline (hexadecimal value `0A`). This happens even for the second string. That is, the `_sjis` introducer does not affect the parser's escape processing.

Example 2:

```
mysql> SET NAMES sjis;
mysql> SELECT HEX('à\n'), HEX(_latin1'à\n');
+------------+-------------------+
| HEX('à\n')  | HEX(_latin1'à\n')  |
+------------+-------------------+
| E05C6E     | E05C6E            |
+------------+-------------------+
```

Here, `character_set_connection` is `sjis`, a character set in which the sequence of `à` followed by `\` (hexadecimal values `05` and `5C`) is a valid multibyte character. Hence, the first two bytes of the string are interpreted as a single `sjis` character, and the `\` is not interpreted as an escape character. The following `n` (hexadecimal value `6E`) is not interpreted as part of an escape sequence. This is true even for the second string; the `_latin1` introducer does not affect escape processing.

10.1.3.7 The National Character Set

Standard SQL defines `NCHAR` or `NATIONAL CHAR` as a way to indicate that a `CHAR` column should use some predefined character set. MySQL uses `utf8` as this predefined character set. For example, these data type declarations are equivalent:

```
CHAR(10) CHARACTER SET utf8
NATIONAL CHARACTER(10)
NCHAR(10)
```

As are these:

```
VARCHAR(10) CHARACTER SET utf8
NATIONAL VARCHAR(10)
NVARCHAR(10)
NCHAR VARCHAR(10)
NATIONAL CHARACTER VARYING(10)
NATIONAL CHAR VARYING(10)
```

You can use `N'literal'` (or `n'literal'`) to create a string in the national character set. These statements are equivalent:

```
SELECT N'some text';
SELECT n'some text';
SELECT _utf8'some text';
```

10.1.3.8 Character Set Introducers

A character string literal, hexadecimal literal, or bit-value literal may have an optional character set introducer and `COLLATE` clause, to designate it as a string that uses a particular character set and collation:

```
[_charset_name] literal [COLLATE collation_name]
```

Character set introducers and the COLLATE clause are implemented according to standard SQL specifications.

Examples:

```
SELECT 'abc';
SELECT _latin1'abc';
SELECT _binary'abc';
SELECT _utf8'abc' COLLATE utf8_danish_ci;

SELECT _latin1 X'4D7953514C';
SELECT _utf8 0x4D7953514C COLLATE utf8_danish_ci;

SELECT _latin1 b'1000001';
SELECT _utf8 0b1000001 COLLATE utf8_danish_ci;
```

The *_charset_name* expression is formally called an *introducer*. It tells the parser, “the string that follows uses character set *charset_name*.” An introducer does not change the string to the introducer character set like CONVERT() would do. It does not change the string value, although padding may occur. The introducer is just a signal.

For character string literals, space between the introducer and the string is permitted but optional.

Character string literals can be designated as binary strings by using the _binary introducer. Hexadecimal literals and bit-value literals are binary strings by default, so _binary is permitted, but normally unnecessary. _binary may be useful to preserve a hexadecimal or bit literal as a binary string in contexts for which the literal is otherwise treated as a number. For example, bit operations permit numeric or binary string arguments in MySQL 8.0 and higher, but treat hexadecimal and bit literals as numbers by default. To explicitly specify binary string context for such literals, use a _binary introducer for at least one of the arguments:

```
mysql> SET @v1 = X'000D' | X'0BC0';
mysql> SET @v2 = _binary X'000D' | X'0BC0';
mysql> SELECT HEX(@v1), HEX(@v2);
+----------+----------+
| HEX(@v1) | HEX(@v2) |
+----------+----------+
| BCD      | 0BCD     |
+----------+----------+
```

The displayed result appears similar for both bit operations, but the result without _binary is a BIGINT value, whereas the result with _binary is a binary string. Due to the difference in result types, the displayed values differ: High-order 0 digits are not displayed for the numeric result.

MySQL determines the character set and collation of a character string literal, hexadecimal literal, or bit-value literal in the following manner:

- If both *_charset_name* and COLLATE *collation_name* are specified, character set *charset_name* and collation *collation_name* are used. *collation_name* must be a permitted collation for *charset_name*.
- If *_charset_name* is specified but COLLATE is not specified, character set *charset_name* and its default collation are used. To see the default collation for each character set, use the SHOW CHARACTER SET statement.
- If *_charset_name* is not specified but COLLATE *collation_name* is specified:
  - For a character string literal, the connection default character set given by the character_set_connection system variable and collation *collation_name* are used. *collation_name* must be a permitted collation for the connection default character set.

  - For a hexadecimal literal or bit-value literal, the only permitted collation is `binary` because these types of literals are binary strings by default.

- Otherwise (neither *`_charset_name`* nor `COLLATE` *`collation_name`* is specified):

  - For a character string literal, the connection default character set and collation given by the `character_set_connection` and `collation_connection` system variables are used.

  - For a hexadecimal literal or bit-value literal, the character set and collation are `binary`.

Examples:

- Nonbinary strings with `latin1` character set and `latin1_german1_ci` collation:

```
SELECT _latin1'Müller' COLLATE latin1_german1_ci;
SELECT _latin1 X'0A0D' COLLATE latin1_german1_ci;
SELECT _latin1 b'0110' COLLATE latin1_german1_ci;
```

- Nonbinary strings with `utf8` character set and its default collation (that is, `utf8_general_ci`):

```
SELECT _utf8'Müller';
SELECT _utf8 X'0A0D';
SELECT _utf8 b'0110';
```

- Binary strings with `binary` character set and its default collation (that is, `binary`):

```
SELECT _binary'Müller';
SELECT X'0A0D';
SELECT b'0110';
```

  The hexadecimal literal and bit-value literal need no introducer because they are binary strings by default.

- A nonbinary string with the connection default character set and `utf8_general_ci` collation (fails if the connection character set is not `utf8`):

```
SELECT 'Müller' COLLATE utf8_general_ci;
```

  This construction (`COLLATE` only) does not work for hexadecimal literals or bit literals because their character set is `binary` no matter the connection character set, and `binary` is not compatible with the `utf8_general_ci` collation. The only permitted `COLLATE` clause in the absence of an introducer is `COLLATE binary`.

- A string with the connection default character set and collation:

```
SELECT 'Müller';
```

For character set literals, an introducer indicates the character set for the following string, but does not change how the parser performs escape processing within the string. Escapes are always interpreted by the parser according to the character set given by `character_set_connection`. For additional discussion and examples, see Section 10.1.3.6, “Character String Literal Character Set and Collation”.

10.1.3.9 Examples of Character Set and Collation Assignment

The following examples show how MySQL determines default character set and collation values.

Example 1: Table and Column Definition

```
```

```
CREATE TABLE t1
(
    c1 CHAR(10) CHARACTER SET latin1 COLLATE latin1_german1_ci
) DEFAULT CHARACTER SET latin2 COLLATE latin2_bin;
```

Here we have a column with a `latin1` character set and a `latin1_german1_ci` collation. The definition is explicit, so that is straightforward. Notice that there is no problem with storing a `latin1` column in a `latin2` table.

Example 2: Table and Column Definition

```
CREATE TABLE t1
(
    c1 CHAR(10) CHARACTER SET latin1
) DEFAULT CHARACTER SET latin1 COLLATE latin1_danish_ci;
```

This time we have a column with a `latin1` character set and a default collation. Although it might seem natural, the default collation is not taken from the table level. Instead, because the default collation for `latin1` is always `latin1_swedish_ci`, column `c1` has a collation of `latin1_swedish_ci` (not `latin1_danish_ci`).

Example 3: Table and Column Definition

```
CREATE TABLE t1
(
    c1 CHAR(10)
) DEFAULT CHARACTER SET latin1 COLLATE latin1_danish_ci;
```

We have a column with a default character set and a default collation. In this circumstance, MySQL checks the table level to determine the column character set and collation. Consequently, the character set for column `c1` is `latin1` and its collation is `latin1_danish_ci`.

Example 4: Database, Table, and Column Definition

```
CREATE DATABASE d1
    DEFAULT CHARACTER SET latin2 COLLATE latin2_czech_ci;
USE d1;
CREATE TABLE t1
(
    c1 CHAR(10)
);
```

We create a column without specifying its character set and collation. We're also not specifying a character set and a collation at the table level. In this circumstance, MySQL checks the database level to determine the table settings, which thereafter become the column settings.) Consequently, the character set for column `c1` is `latin2` and its collation is `latin2_czech_ci`.

10.1.3.10 Compatibility with Other DBMSs

For MaxDB compatibility these two statements are the same:

```
CREATE TABLE t1 (f1 CHAR(N) UNICODE);
CREATE TABLE t1 (f1 CHAR(N) CHARACTER SET ucs2);
```

10.1.4 Connection Character Sets and Collations

Several character set and collation system variables relate to a client's interaction with the server. Some of these have been mentioned in earlier sections:

- The server character set and collation are the values of the `character_set_server` and `collation_server` system variables.

- The character set and collation of the default database are the values of the `character_set_database` and `collation_database` system variables.

Additional character set and collation system variables are involved in handling traffic for the connection between a client and the server. Every client has connection-related character set and collation system variables.

A “connection” is what you make when you connect to the server. The client sends SQL statements, such as queries, over the connection to the server. The server sends responses, such as result sets or error messages, over the connection back to the client. This leads to several questions about character set and collation handling for client connections, each of which can be answered in terms of system variables:

- What character set is the statement in when it leaves the client?

  The server takes the `character_set_client` system variable to be the character set in which statements are sent by the client.

- What character set should the server translate a statement to after receiving it?

  For this, the server uses the `character_set_connection` and `collation_connection` system variables. It converts statements sent by the client from `character_set_client` to `character_set_connection`, xexcept for string literals that have an introducer (for example, `_utf8mb4` or `_latin2`). `collation_connection` is important for comparisons of literal strings. For comparisons of strings with column values, `collation_connection` does not matter because columns have their own collation, which has a higher collation precedence.

- What character set should the server translate to before shipping result sets or error messages back to the client?

  The `character_set_results` system variable indicates the character set in which the server returns query results to the client. This includes result data such as column values, and result metadata such as column names and error messages.

Clients can fine-tune the settings for these variables, or depend on the defaults (in which case, you can skip the rest of this section). If you do not use the defaults, you must change the character settings *for each connection to the server.*

Two statements affect the connection-related character set variables as a group:

- `SET NAMES 'charset_name' [COLLATE 'collation_name']`

  `SET NAMES` indicates what character set the client will use to send SQL statements to the server. Thus, `SET NAMES 'cp1251'` tells the server, “future incoming messages from this client are in character set `cp1251`.” It also specifies the character set that the server should use for sending results back to the client. (For example, it indicates what character set to use for column values if you use a `SELECT` statement.)

  A `SET NAMES 'charset_name'` statement is equivalent to these three statements:

  ```
  SET character_set_client = charset_name;
  SET character_set_results = charset_name;
  SET character_set_connection = charset_name;
  ```

  Setting `character_set_connection` to *`charset_name`* also implicitly sets `collation_connection` to the default collation for *`charset_name`*. It is unnecessary to set that collation explicitly. To specify a particular collation, use the optional `COLLATE` clause:

  ```
  SET NAMES 'charset_name' COLLATE 'collation_name'
  ```

- `SET CHARACTER SET 'charset_name'`

`SET CHARACTER SET` is similar to `SET NAMES` but sets `character_set_connection` and `collation_connection` to `character_set_database` and `collation_database`. A `SET CHARACTER SET *charset_name*` statement is equivalent to these three statements:

```
SET character_set_client = charset_name;
SET character_set_results = charset_name;
SET collation_connection = @@collation_database;
```

Setting `collation_connection` also implicitly sets `character_set_connection` to the character set associated with the collation (equivalent to executing `SET character_set_connection = @@character_set_database`). It is unnecessary to set `character_set_connection` explicitly.

Note

`ucs2`, `utf16`, `utf16le`, and `utf32` cannot be used as a client character set, which means that they do not work for `SET NAMES` or `SET CHARACTER SET`.

The MySQL client programs `mysql`, `mysqladmin`, `mysqlcheck`, `mysqlimport`, and `mysqlshow` determine the default character set to use as follows:

- In the absence of other information, the programs use the compiled-in default character set, usually `utf8mb4` (`latin1` prior to MySQL 8.0.2).
- The programs can autodetect which character set to use based on the operating system setting, such as the value of the `LANG` or `LC_ALL` locale environment variable on Unix systems or the code page setting on Windows systems. For systems on which the locale is available from the OS, the client uses it to set the default character set rather than using the compiled-in default. For example, setting `LANG` to `ru_RU.KOI8-R` causes the `koi8r` character set to be used. Thus, users can configure the locale in their environment for use by MySQL clients.

  The OS character set is mapped to the closest MySQL character set if there is no exact match. If the client does not support the matching character set, it uses the compiled-in default. For example, `ucs2` is not supported as a connection character set. `utf8` and `utf-8` map to `utf8mb4` (to `utf8` prior to MySQL 8.0.2).

  C applications can use character set autodetection based on the OS setting by invoking `mysql_options()` as follows before connecting to the server:

  ```
  mysql_options(mysql,
                MYSQL_SET_CHARSET_NAME,
                MYSQL_AUTODETECT_CHARSET_NAME);
  ```

- The programs support a `--default-character-set` option, which enables users to specify the character set explicitly to override whatever default the client otherwise determines.

When a client connects to the server, it sends the name of the character set to be used for communication with the server. The server uses the name to set the `character_set_client`, `character_set_results`, and `character_set_connection` system variables. In effect, the server performs a `SET NAMES` operation using the character set name.

With the `mysql` client, to use a character set different from the default, you could explicitly execute `SET NAMES` every time you start up. To accomplish the same result more easily, add the `--default-character-set` option setting to your `mysql` command line or in your option file. For example, the following option file setting changes the three connection-related character set variables set to `koi8r` each time you invoke `mysql`:

```
[mysql]
default-character-set=koi8r
```

If you are using the `mysql` client with auto-reconnect enabled (which is not recommended), it is preferable to use the `charset` command rather than `SET NAMES`. For example:

```
mysql> charset utf8
Charset changed
```

The `charset` command issues a `SET NAMES` statement, and also changes the default character set that `mysql` uses when it reconnects after the connection has dropped.

Example: Suppose that `column1` is defined as `CHAR(5) CHARACTER SET latin2`. If you do not say `SET NAMES` or `SET CHARACTER SET`, then for `SELECT column1 FROM t`, the server sends back all the values for `column1` using the character set that the client specified when it connected. On the other hand, if you say `SET NAMES 'latin1'` or `SET CHARACTER SET latin1` before issuing the `SELECT` statement, the server converts the `latin2` values to `latin1` just before sending results back. Conversion may be lossy if there are characters that are not in both character sets.

If you want the server to perform no conversion of result sets or error messages, set `character_set_results` to `NULL` or `binary`:

```
SET character_set_results = NULL;
SET character_set_results = binary;
```

To see the values of the character set and collation system variables that apply to your connection, use these statements:

```
SHOW VARIABLES LIKE 'character_set%';
SHOW VARIABLES LIKE 'collation%';
```

You must also consider the environment within which your MySQL applications execute. See Section 10.1.5, “Configuring Application Character Set and Collation”.

For more information about character sets and error messages, see Section 10.1.6, “Error Message Character Set”.

10.1.5 Configuring Application Character Set and Collation

For applications that store data using the default MySQL character set and collation (`latin1`, `latin1_swedish_ci`), no special configuration should be needed. If applications require data storage using a different character set or collation, you can configure character set information several ways:

- Specify character settings per database. For example, applications that use one database might require `utf8`, whereas applications that use another database might require `sjis`.
- Specify character settings at server startup. This causes the server to use the given settings for all applications that do not make other arrangements.
- Specify character settings at configuration time, if you build MySQL from source. This causes the server to use the given settings for all applications, without having to specify them at server startup.

When different applications require different character settings, the per-database technique provides a good deal of flexibility. If most or all applications use the same character set, specifying character settings at server startup or configuration time may be most convenient.

For the per-database or server-startup techniques, the settings control the character set for data storage. Applications must also tell the server which character set to use for client/server communications, as described in the following instructions.

The examples shown here assume use of the `utf8` character set and `utf8_general_ci` collation.

Specify character settings per database. To create a database such that its tables will use a given default character set and collation for data storage, use a `CREATE DATABASE` statement like this:

```
CREATE DATABASE mydb
  DEFAULT CHARACTER SET utf8
  DEFAULT COLLATE utf8_general_ci;
```

Tables created in the database will use `utf8` and `utf8_general_ci` by default for any character columns.

Applications that use the database should also configure their connection to the server each time they connect. This can be done by executing a `SET NAMES 'utf8'` statement after connecting. The statement can be used regardless of connection method: The `mysql` client, PHP scripts, and so forth.

In some cases, it may be possible to configure the connection to use the desired character set some other way. For example, for connections made using `mysql`, you can specify the `--default-character-set=utf8` command-line option to achieve the same effect as `SET NAMES 'utf8'`.

For more information about configuring client connections, see Section 10.1.4, “Connection Character Sets and Collations”.

If you change the default character set or collation for a database, stored routines that use the database defaults must be dropped and recreated so that they use the new defaults. (In a stored routine, variables with character data types use the database defaults if the character set or collation are not specified explicitly. See Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”.)

Specify character settings at server startup. To select a character set and collation at server startup, use the `--character-set-server` and `--collation-server` options. For example, to specify the options in an option file, include these lines:

```
[mysqld]
character-set-server=utf8
collation-server=utf8_general_ci
```

These settings apply server-wide and apply as the defaults for databases created by any application, and for tables created in those databases.

It is still necessary for applications to configure their connection using `SET NAMES` or equivalent after they connect, as described previously. You might be tempted to start the server with the `--init_connect="SET NAMES 'utf8'"` option to cause `SET NAMES` to be executed automatically for each client that connects. However, this will yield inconsistent results because the `init_connect` value is not executed for users who have the `CONNECTION_ADMIN` or `SUPER` privilege.

Specify character settings at MySQL configuration time. To select a character set and collation when you configure and build MySQL from source, use the `DEFAULT_CHARSET` and `DEFAULT_COLLATION` options for `CMake`:

```
shell> cmake . -DDEFAULT_CHARSET=utf8 \
           -DDEFAULT_COLLATION=utf8_general_ci
```

The resulting server uses `utf8` and `utf8_general_ci` as the default for databases and tables and for client connections. It is unnecessary to use `--character-set-server` and `--collation-server` to specify those defaults at server startup. It is also unnecessary for applications to configure their connection using `SET NAMES` or equivalent after they connect to the server.

Regardless of how you configure the MySQL character set for application use, you must also consider the environment within which those applications execute. If you will send statements using UTF-8 text taken from a file that you create in an editor, you should edit the file with the locale of your environment set to UTF-8 so that the file encoding is correct and so that the operating system handles it correctly. If you use the `mysql` client from within a terminal window, the window must be configured to use UTF-8 or characters may not display properly. For a script that executes in a Web environment, the script must handle character encoding properly for its interaction with the MySQL server, and it must

generate pages that correctly indicate the encoding so that browsers know how to display the content of the pages. For example, you can include this `<meta>` tag within your `<head>` element:

```
<meta http-equiv="Content-Type" content="text/html; charset=utf-8" />
```

10.1.6 Error Message Character Set

This section describes how the MySQL server uses character sets for constructing error messages and returning them to clients. For information about the language of error messages (rather than the character set), see Section 10.2, “Setting the Error Message Language”.

The server constructs error messages using UTF-8 and returns them to clients in the character set specified by the `character_set_results` system variable.

The server constructs error messages as follows:

- The message template uses UTF-8.
- Parameters in the message template are replaced with values that apply to a specific error occurrence:
  - Identifiers such as table or column names use UTF-8 internally so they are copied as is.
  - Character (nonbinary) string values are converted from their character set to UTF-8.
  - Binary string values are copied as is for bytes in the range `0x20` to `0x7E`, and using `\x` hexadecimal encoding for bytes outside that range. For example, if a duplicate-key error occurs for an attempt to insert `0x41CF9F` into a `VARBINARY` unique column, the resulting error message uses UTF-8 with some bytes hexadecimal encoded:

    ```
    Duplicate entry 'A\xC3\x9F' for key 1
    ```

To return a message to the client after it has been constructed, the server converts it from UTF-8 to the character set specified by the `character_set_results` system variable. If `character_set_results` has a value of `NULL` or `binary`, no conversion occurs. No conversion occurs if the variable value is `utf8`, either, because that matches the original error message character set.

For characters that cannot be represented in `character_set_results`, some encoding may occur during the conversion. The encoding uses Unicode code point values:

- Characters in the Basic Multilingual Plane (BMP) range (`0x0000` to `0xFFFF`) are written using `\`*`nnnn`* notation.
- Characters outside the BMP range (`0x01000` to `0x10FFFF`) are written using `\+`*`nnnnnn`* notation.

Clients can set `character_set_results` to control the character set in which they receive error messages. The variable can be set directly, or indirectly by means such as `SET NAMES`. For more information about `character_set_results`, see Section 10.1.4, “Connection Character Sets and Collations”.

The encoding that occurs during the conversion to `character_set_results` before returning error messages to clients can result in different message content compared to earlier versions (before MySQL 5.5). For example, if an error occurs for an attempt to drop a table named ㌀ (KATAKANA LETTER PE) and `character_set_results` is a character set such as `latin1` that does not contain that character, the resulting message sent to the client has an encoded table name:

```
ERROR 1051 (42S02): Unknown table '\30DA'
```

Before MySQL 5.5, the name is not encoded:

```
ERROR 1051 (42S02): Unknown table 'ペ'
```

10.1.7 Column Character Set Conversion

To convert a binary or nonbinary string column to use a particular character set, use `ALTER TABLE`. For successful conversion to occur, one of the following conditions must apply:

- If the column has a binary data type (`BINARY`, `VARBINARY`, `BLOB`), all the values that it contains must be encoded using a single character set (the character set you're converting the column to). If you use a binary column to store information in multiple character sets, MySQL has no way to know which values use which character set and cannot convert the data properly.
- If the column has a nonbinary data type (`CHAR`, `VARCHAR`, `TEXT`), its contents should be encoded in the column character set, not some other character set. If the contents are encoded in a different character set, you can convert the column to use a binary data type first, and then to a nonbinary column with the desired character set.

Suppose that a table `t` has a binary column named `col1` defined as `VARBINARY(50)`. Assuming that the information in the column is encoded using a single character set, you can convert it to a nonbinary column that has that character set. For example, if `col1` contains binary data representing characters in the `greek` character set, you can convert it as follows:

```
ALTER TABLE t MODIFY col1 VARCHAR(50) CHARACTER SET greek;
```

If your original column has a type of `BINARY(50)`, you could convert it to `CHAR(50)`, but the resulting values will be padded with `0x00` bytes at the end, which may be undesirable. To remove these bytes, use the `TRIM()` function:

```
UPDATE t SET col1 = TRIM(TRAILING 0x00 FROM col1);
```

Suppose that table `t` has a nonbinary column named `col1` defined as `CHAR(50) CHARACTER SET latin1` but you want to convert it to use `utf8` so that you can store values from many languages. The following statement accomplishes this:

```
ALTER TABLE t MODIFY col1 CHAR(50) CHARACTER SET utf8;
```

Conversion may be lossy if the column contains characters that are not in both character sets.

A special case occurs if you have old tables from before MySQL 4.1 where a nonbinary column contains values that actually are encoded in a character set different from the server's default character set. For example, an application might have stored `sjis` values in a column, even though MySQL's default character set was `latin1`. It is possible to convert the column to use the proper character set but an additional step is required. Suppose that the server's default character set was `latin1` and `col1` is defined as `CHAR(50)` but its contents are `sjis` values. The first step is to convert the column to a binary data type, which removes the existing character set information without performing any character conversion:

```
ALTER TABLE t MODIFY col1 BLOB;
```

The next step is to convert the column to a nonbinary data type with the proper character set:

```
ALTER TABLE t MODIFY col1 CHAR(50) CHARACTER SET sjis;
```

This procedure requires that the table not have been modified already with statements such as `INSERT` or `UPDATE` after an upgrade to MySQL 4.1 or later. In that case, MySQL would store new values in the column using `latin1`, and the column will contain a mix of `sjis` and `latin1` values and cannot be converted properly.

If you specified attributes when creating a column initially, you should also specify them when altering the table with `ALTER TABLE`. For example, if you specified `NOT NULL` and an explicit `DEFAULT` value, you should also provide them in the `ALTER TABLE` statement. Otherwise, the resulting column definition will not include those attributes.

To convert all character columns in a table, the `ALTER TABLE ... CONVERT TO CHARACTER SET charset` statement may be useful. See Section 13.1.7, “ALTER TABLE Syntax”.

10.1.8 Collation Issues

The following sections discuss various aspects of character set collations.

10.1.8.1 Using COLLATE in SQL Statements

With the `COLLATE` clause, you can override whatever the default collation is for a comparison. `COLLATE` may be used in various parts of SQL statements. Here are some examples:

- With `ORDER BY`:

```
SELECT k
FROM t1
ORDER BY k COLLATE latin1_german2_ci;
```

- With `AS`:

```
SELECT k COLLATE latin1_german2_ci AS k1
FROM t1
ORDER BY k1;
```

- With `GROUP BY`:

```
SELECT k
FROM t1
GROUP BY k COLLATE latin1_german2_ci;
```

- With aggregate functions:

```
SELECT MAX(k COLLATE latin1_german2_ci)
FROM t1;
```

- With `DISTINCT`:

```
SELECT DISTINCT k COLLATE latin1_german2_ci
FROM t1;
```

- With `WHERE`:

```
    SELECT *
    FROM t1
    WHERE _latin1 'Müller' COLLATE latin1_german2_ci = k;
```

```
    SELECT *
    FROM t1
    WHERE k LIKE _latin1 'Müller' COLLATE latin1_german2_ci;
```

- With `HAVING`:

```
SELECT k
FROM t1
```

```
GROUP BY k
HAVING k = _latin1 'Müller' COLLATE latin1_german2_ci;
```

10.1.8.2 COLLATE Clause Precedence

The `COLLATE` clause has high precedence (higher than `||`), so the following two expressions are equivalent:

```
x || y COLLATE z
x || (y COLLATE z)
```

10.1.8.3 Character Set and Collation Compatibility

Each character set has one or more collations, but each collation is associated with one and only one character set. Therefore, the following statement causes an error message because the `latin2_bin` collation is not legal with the `latin1` character set:

```
mysql> SELECT _latin1 'x' COLLATE latin2_bin;
ERROR 1253 (42000): COLLATION 'latin2_bin' is not valid
for CHARACTER SET 'latin1'
```

10.1.8.4 Collation Coercibility in Expressions

In the great majority of statements, it is obvious what collation MySQL uses to resolve a comparison operation. For example, in the following cases, it should be clear that the collation is the collation of column `x`:

```
SELECT x FROM T ORDER BY x;
SELECT x FROM T WHERE x = x;
SELECT DISTINCT x FROM T;
```

However, with multiple operands, there can be ambiguity. For example:

```
SELECT x FROM T WHERE x = 'Y';
```

Should the comparison use the collation of the column `x`, or of the string literal `'Y'`? Both `x` and `'Y'` have collations, so which collation takes precedence?

A mix of collations may also occur in contexts other than comparison. For example, a multiple-argument concatenation operation such as `CONCAT(x,'Y')` combines its arguments to produce a single string. What collation should the result have?

To resolve questions like these, MySQL checks whether the collation of one item can be coerced to the collation of the other. MySQL assigns coercibility values as follows:

- An explicit `COLLATE` clause has a coercibility of 0 (not coercible at all).
- The concatenation of two strings with different collations has a coercibility of 1.
- The collation of a column or a stored routine parameter or local variable has a coercibility of 2.
- A “system constant” (the string returned by functions such as `USER()` or `VERSION()`) has a coercibility of 3.
- The collation of a literal has a coercibility of 4.
- The collation of a numeric or temporal value has a coercibility of 5.
- `NULL` or an expression that is derived from `NULL` has a coercibility of 6.

MySQL uses coercibility values with the following rules to resolve ambiguities:

- Use the collation with the lowest coercibility value.
- If both sides have the same coercibility, then:
  - If both sides are Unicode, or both sides are not Unicode, it is an error.
  - If one of the sides has a Unicode character set, and another side has a non-Unicode character set, the side with Unicode character set wins, and automatic character set conversion is applied to the non-Unicode side. For example, the following statement does not return an error:

    ```
    SELECT CONCAT(utf8_column, latin1_column) FROM t1;
    ```

    It returns a result that has a character set of `utf8` and the same collation as `utf8_column`. Values of `latin1_column` are automatically converted to `utf8` before concatenating.
  - For an operation with operands from the same character set but that mix a `_bin` collation and a `_ci` or `_cs` collation, the `_bin` collation is used. This is similar to how operations that mix nonbinary and binary strings evaluate the operands as binary strings, except that it is for collations rather than data types.

Although automatic conversion is not in the SQL standard, the standard does say that every character set is (in terms of supported characters) a “subset” of Unicode. Because it is a well-known principle that “what applies to a superset can apply to a subset,” we believe that a collation for Unicode can apply for comparisons with non-Unicode strings.

The following table illustrates some applications of the preceding rules.

Comparison	Collation Used
`column1 = 'A'`	Use collation of `column1`
`column1 = 'A' COLLATE x`	Use collation of `'A' COLLATE x`
`column1 COLLATE x = 'A' COLLATE y`	Error

To determine the coercibility of a string expression, use the `COERCIBILITY()` function (see Section 12.14, “Information Functions”):

```
mysql> SELECT COERCIBILITY('A' COLLATE latin1_swedish_ci);
        -> 0
mysql> SELECT COERCIBILITY(VERSION());
        -> 3
mysql> SELECT COERCIBILITY('A');
        -> 4
mysql> SELECT COERCIBILITY(1000);
        -> 5
```

For implicit conversion of a numeric or temporal value to a string, such as occurs for the argument `1` in the expression `CONCAT(1, 'abc')`, the result is a character (nonbinary) string that has a character set and collation determined by the `character_set_connection` and `collation_connection` system variables. See Section 12.2, “Type Conversion in Expression Evaluation”.

10.1.8.5 The binary Collation Compared to _bin Collations

This section describes how the `binary` collation for binary strings compares to the `_bin` collations for nonbinary strings.

Binary strings (as stored using the `BINARY`, `VARBINARY`, and `BLOB` data types) have a character set and collation named `binary`. Binary strings are sequences of bytes and the numeric values of those bytes determine comparison and sort order.

Nonbinary strings (as stored using the `CHAR`, `VARCHAR`, and `TEXT` data types) have a character set and collation other than `binary`. A given nonbinary character set can have several collations, each of which defines a particular comparison and sort order for the characters in the set. One of these is the binary collation for the character set, indicated by a `_bin` suffix in the collation name. For example, the binary collations for `latin1` and `utf8` are named `latin1_bin` and `utf8_bin`, respectively.

The `binary` collation differs from the `_bin` collations in several respects.

The unit for comparison and sorting. Binary strings are sequences of bytes. For the `binary` collation, comparison and sorting are based on numeric byte values. Nonbinary strings are sequences of characters, which might be multibyte. Collations for nonbinary strings define an ordering of the character values for comparison and sorting. For the `_bin` collation, this ordering is based on numeric character code values, which is similar to ordering for binary strings except that character code values might be multibyte.

Character set conversion. A nonbinary string has a character set and is automatically converted to another character set in many cases, even when the string has a `_bin` collation:

- When assigning column values from another column that has a different character set:

```
UPDATE t1 SET utf8_bin_column=latin1_column;
INSERT INTO t1 (latin1_column) SELECT utf8_bin_column FROM t2;
```

- When assigning column values for `INSERT` or `UPDATE` using a string literal:

```
SET NAMES latin1;
INSERT INTO t1 (utf8_bin_column) VALUES ('string-in-latin1');
```

- When sending results from the server to a client:

```
SET NAMES latin1;
SELECT utf8_bin_column FROM t2;
```

For binary string columns, no conversion occurs. For the preceding cases, the string value is copied byte-wise.

Lettercase conversion. Collations for nonbinary character sets provide information about lettercase of characters, so characters in a nonbinary string can be converted from one lettercase to another, even for `_bin` collations that ignore lettercase for ordering:

```
mysql> SET NAMES latin1 COLLATE latin1_bin;
mysql> SELECT LOWER('aA'), UPPER('zZ');
+-------------+-------------+
| LOWER('aA') | UPPER('zZ') |
+-------------+-------------+
| aa          | ZZ          |
+-------------+-------------+
```

The concept of lettercase does not apply to bytes in a binary string. To perform lettercase conversion, the string must be converted to a nonbinary string:

```
mysql> SET NAMES binary;
mysql> SELECT LOWER('aA'), LOWER(CONVERT('aA' USING latin1));
+-------------+-----------------------------------+
| LOWER('aA') | LOWER(CONVERT('aA' USING latin1)) |
+-------------+-----------------------------------+
| aA          | aa                                |
+-------------+-----------------------------------+
```

Trailing space handling in comparisons.

Most MySQL collations have a pad attribute of PAD SPACE. The exceptions are Unicode collations based on UCA 9.0.0 and higher, which have a pad attribute of NO PAD. (see Section 10.1.10.1, “Unicode Character Sets”).

To determine the pad attribute for a collation, use the `INFORMATION_SCHEMA COLLATIONS` table, which has a `PAD_ATTRIBUTE` column.

The pad attribute determines how trailing spaces are treated for comparison of nonbinary strings (`CHAR`, `VARCHAR`, and `TEXT` values). NO PAD collations treat spaces at the end of strings like any other character. For PAD SPACE collations, trailing spaces are insignificant in comparisons; strings are compared without regard to any trailing spaces:

```
mysql> SET NAMES utf8 COLLATE utf8_bin;
mysql> SELECT 'a ' = 'a';
+------------+
| 'a ' = 'a' |
+------------+
|          1 |
+------------+
```

For binary strings, all characters are significant in comparisons, including trailing spaces:

```
mysql> SET NAMES binary;
mysql> SELECT 'a ' = 'a';
+------------+
| 'a ' = 'a' |
+------------+
|          0 |
+------------+
```

Trailing space handling for inserts and retrievals. `CHAR(N)` columns store nonbinary strings. Values shorter than *`N`* characters are extended with spaces on insertion. For retrieval, trailing spaces are removed.

`BINARY(N)` columns store binary strings. Values shorter than *`N`* bytes are extended with `0x00` bytes on insertion. For retrieval, nothing is removed; a value of the declared length is always returned.

```
mysql> CREATE TABLE t1 (
         a CHAR(10) CHARACTER SET utf8 COLLATE utf8_bin,
         b BINARY(10)
       );
mysql> INSERT INTO t1 VALUES ('a','a');
mysql> SELECT HEX(a), HEX(b) FROM t1;
+--------+----------------------+
| HEX(a) | HEX(b)               |
+--------+----------------------+
| 61     | 61000000000000000000 |
+--------+----------------------+
```

10.1.8.6 Examples of the Effect of Collation

Example 1: Sorting German Umlauts

Suppose that column `X` in table `T` has these `latin1` column values:

```
Muffler
Müller
MX Systems
MySQL
```

Suppose also that the column values are retrieved using the following statement:

```
SELECT X FROM T ORDER BY X COLLATE collation_name;
```

The following table shows the resulting order of the values if we use `ORDER BY` with different collations.

latin1_swedish_ci	latin1_german1_ci	latin1_german2_ci
Muffler	Muffler	Müller
MX Systems	Müller	Muffler
Müller	MX Systems	MX Systems
MySQL	MySQL	MySQL

The character that causes the different sort orders in this example is the U with two dots over it (ü), which the Germans call “U-umlaut.”

- The first column shows the result of the `SELECT` using the Swedish/Finnish collating rule, which says that U-umlaut sorts with Y.
- The second column shows the result of the `SELECT` using the German DIN-1 rule, which says that U-umlaut sorts with U.
- The third column shows the result of the `SELECT` using the German DIN-2 rule, which says that U-umlaut sorts with UE.

Example 2: Searching for German Umlauts

Suppose that you have three tables that differ only by the character set and collation used:

```
mysql> SET NAMES utf8;
mysql> CREATE TABLE german1 (
         c CHAR(10)
       ) CHARACTER SET latin1 COLLATE latin1_german1_ci;
mysql> CREATE TABLE german2 (
         c CHAR(10)
       ) CHARACTER SET latin1 COLLATE latin1_german2_ci;
mysql> CREATE TABLE germanutf8 (
         c CHAR(10)
       ) CHARACTER SET utf8 COLLATE utf8_unicode_ci;
```

Each table contains two records:

```
mysql> INSERT INTO german1 VALUES ('Bar'), ('Bär');
mysql> INSERT INTO german2 VALUES ('Bar'), ('Bär');
mysql> INSERT INTO germanutf8 VALUES ('Bar'), ('Bär');
```

Two of the above collations have an `A = Ä` equality, and one has no such equality (`latin1_german2_ci`). For that reason, you'll get these results in comparisons:

```
mysql> SELECT * FROM german1 WHERE c = 'Bär';
+------+
| c    |
+------+
| Bar  |
| Bär  |
+------+
mysql> SELECT * FROM german2 WHERE c = 'Bär';
+------+
| c    |
+------+
| Bär  |
+------+
mysql> SELECT * FROM germanutf8 WHERE c = 'Bär';
+------+
| c    |
+------+
| Bar  |
```

```
| Bär  |
+------+
```

This is not a bug but rather a consequence of the sorting properties of `latin1_german1_ci` and `utf8_unicode_ci` (the sorting shown is done according to the German DIN 5007 standard).

10.1.8.7 Using Collation in INFORMATION_SCHEMA Searches

String columns in `INFORMATION_SCHEMA` tables have a collation of `utf8_general_ci`, which is case insensitive. However, for values that correspond to objects that are represented in the file system, such as databases and tables, searches in `INFORMATION_SCHEMA` string columns can be case sensitive or insensitive, depending on the characteristics of the underlying file system and the value of the `lower_case_table_names` system variable. For example, searches may be case sensitive if the file system is case sensitive. This section describes this behavior and how to modify it if necessary.

Suppose that a query searches the `SCHEMATA.SCHEMA_NAME` column for the `test` database. On Linux, file systems are case sensitive, so comparisons of `SCHEMATA.SCHEMA_NAME` with `'test'` match, but comparisons with `'TEST'` do not:

```
mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'test';
+-------------+
| SCHEMA_NAME |
+-------------+
| test        |
+-------------+

mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'TEST';
Empty set (0.00 sec)
```

These results occur with the `lower_case_table_names` system system variable set to 0. Changing the value of `lower_case_table_names` to 1 or 2 causes the second query to return the same (nonempty) result as the first query.

On Windows or OS X, file systems are not case sensitive, so comparisons match both `'test'` and `'TEST'`:

```
mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'test';
+-------------+
| SCHEMA_NAME |
+-------------+
| test        |
+-------------+

mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'TEST';
+-------------+
| SCHEMA_NAME |
+-------------+
| TEST        |
+-------------+
```

The value of `lower_case_table_names` makes no difference in this context.

The preceding behavior occurs because the `utf8_general_ci` collation is not used for `INFORMATION_SCHEMA` queries when searching for values that correspond to objects represented in the file system.

If the result of a string operation on an `INFORMATION_SCHEMA` column differs from expectations, a workaround is to use an explicit `COLLATE` clause to force a suitable collation (see Section 10.1.8.1, “Using COLLATE in SQL Statements”). For example, to perform a case-insensitive search, use `COLLATE` with the `INFORMATION_SCHEMA` column name:

```
mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME COLLATE utf8_general_ci = 'test';
+-------------+
| SCHEMA_NAME |
+-------------+
| test        |
+-------------+

mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME COLLATE utf8_general_ci = 'TEST';
+-------------+
| SCHEMA_NAME |
+-------------+
| test        |
+-------------+
```

You can also use the `UPPER()` or `LOWER()` function:

```
WHERE UPPER(SCHEMA_NAME) = 'TEST'
WHERE LOWER(SCHEMA_NAME) = 'test'
```

Although a case-insensitive comparison can be performed even on platforms with case-sensitive file systems, as just shown, it is not necessarily always the right thing to do. On such platforms, it is possible to have multiple objects with names that differ only in lettercase. For example, tables named `city`, `CITY`, and `City` can all exist simultaneously. Consider whether a search should match all such names or just one and write queries accordingly. The first of the following comparisons (with `utf8_bin`) is case sensitive; the others are not:

```
WHERE TABLE_NAME COLLATE utf8_bin = 'City'
WHERE TABLE_NAME COLLATE utf8_general_ci = 'city'
WHERE UPPER(TABLE_NAME) = 'CITY'
WHERE LOWER(TABLE_NAME) = 'city'
```

Searches in `INFORMATION_SCHEMA` string columns for values that refer to `INFORMATION_SCHEMA` itself do use the `utf8_general_ci` collation because `INFORMATION_SCHEMA` is a “virtual” database not represented in the file system. For example, comparisons with `SCHEMATA.SCHEMA_NAME` match `'information_schema'` or `'INFORMATION_SCHEMA'` regardless of platform:

```
mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'information_schema';
+--------------------+
| SCHEMA_NAME        |
+--------------------+
| information_schema |
+--------------------+

mysql> SELECT SCHEMA_NAME FROM INFORMATION_SCHEMA.SCHEMATA
       WHERE SCHEMA_NAME = 'INFORMATION_SCHEMA';
+--------------------+
| SCHEMA_NAME        |
+--------------------+
| information_schema |
+--------------------+
```

10.1.9 Unicode Support

The Unicode Standard includes characters from the Basic Multilingual Plane (BMP) and supplementary characters that lie outside the BMP. This section describes support for Unicode in MySQL. For information about the Unicode Standard itself, visit the Unicode Consortium Web site.

BMP characters have these characteristics:

- Their code point values are between 0 and 65535 (or `U+0000` and `U+FFFF`).

- They can be encoded in a variable-length encoding using 8, 16, or 24 bits (1 to 3 bytes).
- They can be encoded in a fixed-length encoding using 16 bits (2 bytes).
- They are sufficient for almost all characters in major languages.

Supplementary characters lie outside the BMP. Their code point values are between `U+10000` and `U+10FFFF`). Unicode support for supplementary characters requires character sets that have a range outside BMP characters and therefore take more space than BMP characters.

MySQL supports these Unicode character sets:

- `utf8`, a UTF-8 encoding of the Unicode character set using one to three bytes per character.
- `ucs2`, the UCS-2 encoding of the Unicode character set using two bytes per character.
- `utf8mb4`, a UTF-8 encoding of the Unicode character set using one to four bytes per character.
- `utf16`, the UTF-16 encoding for the Unicode character set using two or four bytes per character. Like `ucs2` but with an extension for supplementary characters.
- `utf16le`, the UTF-16LE encoding for the Unicode character set. Like `utf16` but little-endian rather than big-endian.
- `utf32`, the UTF-32 encoding for the Unicode character set using four bytes per character.

Table 10.2, “Unicode Character Set General Characteristics”, summarizes the general characteristics of Unicode character sets supported by MySQL.

Table 10.2 Unicode Character Set General Characteristics

Character Set	Supported Characters	Required Storage Per Character
`utf8`	BMP only	1, 2, or 3 bytes
`ucs2`	BMP only	2 bytes
`utf8mb4`	BMP and supplementary	1, 2, 3, or 4 bytes
`utf16`	BMP and supplementary	2 or 4 bytes
`utf16le`	BMP and supplementary	2 or 4 bytes
`utf32`	BMP and supplementary	4 bytes

Characters outside the BMP compare as REPLACEMENT CHARACTER and convert to `'?'` when converted to a Unicode character set that supports only BMP characters (`utf8` or `ucs2`).

If you use character sets that support supplementary characters and thus are “wider” than the BMP-only `utf8` and `ucs2` character sets, there are potential incompatibility issues for your applications; see Section 10.1.9.8, “Converting Between 3-Byte and 4-Byte Unicode Character Sets”. That section also describes how to convert tables from `utf8` to the (4-byte) `utf8mb4` character set, and what constraints may apply in doing so.

A similar set of collations is available for most Unicode character sets. For example, each has a Danish collation, the names of which are `ucs2_danish_ci`, `utf16_danish_ci`, `utf32_danish_ci`, `utf8_danish_ci`, and `utf8mb4_danish_ci`. The exception is `utf16le`, which has only two collations. For information about Unicode collations and their differentiating properties, including collation properties for supplementary characters, see Section 10.1.10.1, “Unicode Character Sets”.

The MySQL implementation of UCS-2, UTF-16, and UTF-32 stores characters in big-endian byte order and does not use a byte order mark (BOM) at the beginning of values. Other database systems might use little-endian byte order or a BOM. In such cases, conversion of values will need to be performed when transferring data between those systems and MySQL. The implementation of UTF-16LE is little-endian.

MySQL uses no BOM for UTF-8 values.

Client applications that communicate with the server using Unicode should set the client character set accordingly; for example, by issuing a `SET NAMES 'utf8'` statement. `ucs2`, `utf16`, `utf16le`, and `utf32` cannot be used as a client character set, which means that they do not work for `SET NAMES` or `SET CHARACTER SET`. (See Section 10.1.4, “Connection Character Sets and Collations”.)

The following sections provide additional detail on the Unicode character sets in MySQL.

10.1.9.1 The utf8 Character Set (3-Byte UTF-8 Unicode Encoding)

UTF-8 (Unicode Transformation Format with 8-bit units) is an alternative way to store Unicode data. It is implemented according to RFC 3629, which describes encoding sequences that take from one to four bytes. (An older standard for UTF-8 encoding, RFC 2279, describes UTF-8 sequences that take from one to six bytes. RFC 3629 renders RFC 2279 obsolete; for this reason, sequences with five and six bytes are no longer used.)

The idea of UTF-8 is that various Unicode characters are encoded using byte sequences of different lengths:

- Basic Latin letters, digits, and punctuation signs use one byte.
- Most European and Middle East script letters fit into a 2-byte sequence: extended Latin letters (with tilde, macron, acute, grave and other accents), Cyrillic, Greek, Armenian, Hebrew, Arabic, Syriac, and others.
- Korean, Chinese, and Japanese ideographs use 3-byte or 4-byte sequences.

The `utf8` character set in MySQL has these characteristics:

- No support for supplementary characters (BMP characters only).
- A maximum of three bytes per multibyte character.

Exactly the same set of characters is available in `utf8` and `ucs2`. That is, they have the same repertoire.

Tip

To save space with UTF-8, use `VARCHAR` instead of `CHAR`. Otherwise, MySQL must reserve three bytes for each character in a `CHAR CHARACTER SET utf8` column because that is the maximum possible character length. For example, MySQL must reserve 30 bytes for a `CHAR(10) CHARACTER SET utf8` column.

For additional information about data type storage, see Section 11.8, “Data Type Storage Requirements”. For information about `InnoDB` physical row storage, including how `InnoDB` tables that use `COMPACT` row format handle UTF-8 `CHAR(N)` columns internally, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

10.1.9.2 The utf8mb3 Character Set (Alias for utf8)

The `utf8` character set uses a maximum of three bytes per character. To make this character limit explicit (much as the limit of four bytes per character is explicit in the `utf8mb4` character set name, use the character set name `utf8mb3`, which is an alias for `utf8`. `utf8mb3` can be used in `CHARACTER SET` clauses, and `utf8mb3_collation_substring` in `COLLATE` clauses, where `collation_substring` is `bin`, `czech_ci`, `danish_ci`, `esperanto_ci`, `estonian_ci`, and so forth. For example:

```
CREATE TABLE t (s1 CHAR(1) CHARACTER SET utf8mb3;
```

```
SELECT * FROM t WHERE s1 COLLATE utf8mb3_general_ci = 'x';
DECLARE x VARCHAR(5) CHARACTER SET utf8mb3 COLLATE utf8mb3_danish_ci;
SELECT CAST('a' AS CHAR CHARACTER SET utf8) COLLATE utf8_czech_ci;
```

MySQL immediately converts instances of `utf8mb3` in an alias to `utf8`, so in statements such as `SHOW CREATE TABLE` or `SELECT CHARACTER_SET_NAME FROM INFORMATION_SCHEMA.COLUMNS` or `SELECT COLLATION_NAME FROM INFORMATION_SCHEMA.COLUMNS`, users will see the true name, `utf8` or `utf8_collation_substring`.

The `utf8mb3` alias is also valid in certain places other than `CHARACTER SET` clauses. For example, these are legal:

```
mysqld --character-set-server=utf8mb3
SET NAMES 'utf8mb3'; /* and other SET statements that have similar effect */
SELECT _utf8mb3 'a';
```

There is no `utf8mb3` alias for the corresponding `utf8` collation for collation names that include a version number to indicate the Unicode Collation Algorithm version on which the collation is based (for example, `utf8_unicode_520_ci`).

10.1.9.3 The utf8mb4 Character Set (4-Byte UTF-8 Unicode Encoding)

The character set named `utf8` uses a maximum of three bytes per character and contains only BMP characters. The `utf8mb4` character set uses a maximum of four bytes per character and supports supplementary characters:

- For a BMP character, `utf8` and `utf8mb4` have identical storage characteristics: same code values, same encoding, same length.
- For a supplementary character, `utf8` cannot store the character at all, whereas `utf8mb4` requires four bytes to store it. Because `utf8` cannot store the character at all, you have no supplementary characters in `utf8` columns and need not worry about converting characters or losing data when upgrading `utf8` data from older versions of MySQL.

`utf8mb4` is a superset of `utf8`, so for an operation such as the following concatenation, the result has character set `utf8mb4` and the collation of `utf8mb4_col`:

```
SELECT CONCAT(utf8_col, utf8mb4_col);
```

Similarly, the following comparison in the `WHERE` clause works according to the collation of `utf8mb4_col`:

```
SELECT * FROM utf8_tbl, utf8mb4_tbl
WHERE utf8_tbl.utf8_col = utf8mb4_tbl.utf8mb4_col;
```

Tip: To save space with `utf8mb4`, use `VARCHAR` instead of `CHAR`. Otherwise, MySQL must reserve four bytes for each character in a `CHAR CHARACTER SET utf8mb4` column because that is the maximum possible length. For example, MySQL must reserve 40 bytes for a `CHAR(10) CHARACTER SET utf8mb4` column.

10.1.9.4 The ucs2 Character Set (UCS-2 Unicode Encoding)

In UCS-2, every character is represented by a 2-byte Unicode code with the most significant byte first. For example: `LATIN CAPITAL LETTER A` has the code `0x0041` and it is stored as a 2-byte sequence: `0x00 0x41`. `CYRILLIC SMALL LETTER YERU` (Unicode `0x044B`) is stored as a 2-byte sequence: `0x04 0x4B`. For Unicode characters and their codes, please refer to the Unicode Consortium Web site.

In MySQL, the `ucs2` character set is a fixed-length 16-bit encoding for Unicode BMP characters.

10.1.9.5 The utf16 Character Set (UTF-16 Unicode Encoding)

The `utf16` character set is the `ucs2` character set with an extension that enables encoding of supplementary characters:

- For a BMP character, `utf16` and `ucs2` have identical storage characteristics: same code values, same encoding, same length.
- For a supplementary character, `utf16` has a special sequence for representing the character using 32 bits. This is called the “surrogate” mechanism: For a number greater than `0xffff`, take 10 bits and add them to `0xd800` and put them in the first 16-bit word, take 10 more bits and add them to `0xdc00` and put them in the next 16-bit word. Consequently, all supplementary characters require 32 bits, where the first 16 bits are a number between `0xd800` and `0xdbff`, and the last 16 bits are a number between `0xdc00` and `0xdfff`. Examples are in Section 15.5 Surrogates Area of the Unicode 4.0 document.

Because `utf16` supports surrogates and `ucs2` does not, there is a validity check that applies only in `utf16`: You cannot insert a top surrogate without a bottom surrogate, or vice versa. For example:

```
INSERT INTO t (ucs2_column) VALUES (0xd800); /* legal */
INSERT INTO t (utf16_column)VALUES (0xd800); /* illegal */
```

There is no validity check for characters that are technically valid but are not true Unicode (that is, characters that Unicode considers to be “unassigned code points” or “private use” characters or even “illegals” like `0xffff`). For example, since `U+F8FF` is the Apple Logo, this is legal:

```
INSERT INTO t (utf16_column)VALUES (0xf8ff); /* legal */
```

Such characters cannot be expected to mean the same thing to everyone.

Because MySQL must allow for the worst case (that one character requires four bytes) the maximum length of a `utf16` column or index is only half of the maximum length for a `ucs2` column or index. For example, the maximum length of a `MEMORY` table index key is 3072 bytes, so these statements create tables with the longest permitted indexes for `ucs2` and `utf16` columns:

```
CREATE TABLE tf (s1 VARCHAR(1536) CHARACTER SET ucs2) ENGINE=MEMORY;
CREATE INDEX i ON tf (s1);
CREATE TABLE tg (s1 VARCHAR(768) CHARACTER SET utf16) ENGINE=MEMORY;
CREATE INDEX i ON tg (s1);
```

10.1.9.6 The utf16le Character Set (UTF-16LE Unicode Encoding)

This is the same as `utf16` but is little-endian rather than big-endian.

10.1.9.7 The utf32 Character Set (UTF-32 Unicode Encoding)

The `utf32` character set is fixed length (like `ucs2` and unlike `utf16`). `utf32` uses 32 bits for every character, unlike `ucs2` (which uses 16 bits for every character), and unlike `utf16` (which uses 16 bits for some characters and 32 bits for others).

`utf32` takes twice as much space as `ucs2` and more space than `utf16`, but `utf32` has the same advantage as `ucs2` that it is predictable for storage: The required number of bytes for `utf32` equals the number of characters times 4. Also, unlike `utf16`, there are no tricks for encoding in `utf32`, so the stored value equals the code value.

To demonstrate how the latter advantage is useful, here is an example that shows how to determine a `utf8mb4` value given the `utf32` code value:

```
/* Assume code value = 100cc LINEAR B WHEELED CHARIOT */
```

```
CREATE TABLE tmp (utf32_col CHAR(1) CHARACTER SET utf32,
                  utf8mb4_col CHAR(1) CHARACTER SET utf8mb4);
INSERT INTO tmp VALUES (0x000100cc,NULL);
UPDATE tmp SET utf8mb4_col = utf32_col;
SELECT HEX(utf32_col),HEX(utf8mb4_col) FROM tmp;
```

MySQL is very forgiving about additions of unassigned Unicode characters or private-use-area characters. There is in fact only one validity check for `utf32`: No code value may be greater than `0x10ffff`. For example, this is illegal:

```
INSERT INTO t (utf32_column) VALUES (0x110000); /* illegal */
```

10.1.9.8 Converting Between 3-Byte and 4-Byte Unicode Character Sets

This section describes issues that you may face when converting from the `utf8` character set to the `utf8mb4` character set, or vice versa.

Note

The discussion here focuses primarily on converting between `utf8` and `utf8mb4`, but similar principles apply to converting between the `ucs2` character set and character sets such as `utf16` or `utf32`.

The `utf8` and `utf8mb4` character sets differ as follows:

- `utf8` supports only characters in the Basic Multilingual Plane (BMP). `utf8mb4` additionally supports supplementary characters that lie outside the BMP.
- `utf8` uses a maximum of three bytes per character. `utf8mb4` uses a maximum of four bytes per character.

One advantage of converting from `ut8` to `utf8mb4` is that this enables applications to use supplementary characters. One tradeoff is that this may increase data storage space requirements.

In most respects, converting from `utf8` to `utf8mb4` should present few problems. These are the primary potential areas of incompatibility:

- For the variable-length character data types (`VARCHAR` and the `TEXT` types), the maximum permitted length in characters is less for `utf8mb4` columns than for `utf8` columns.
- For all character data types (`CHAR`, `VARCHAR`, and the `TEXT` types), the maximum number of characters that can be indexed is less for `utf8mb4` columns than for `utf8` columns.

Consequently, to convert tables from `utf8` to `utf8mb4`, it may be necessary to change some column or index definitions.

Tables can be converted from `utf8` to `utf8mb4` by using `ALTER TABLE`. Suppose that a table was originally defined as follows:

```
CREATE TABLE t1 (
  col1 CHAR(10) CHARACTER SET utf8 COLLATE utf8_unicode_ci NOT NULL,
  col2 CHAR(10) CHARACTER SET utf8 COLLATE utf8_bin NOT NULL
) CHARACTER SET utf8;
```

The following statement converts `t1` to use `utf8mb4`:

```
ALTER TABLE t1
  DEFAULT CHARACTER SET utf8mb4,
  MODIFY col1 CHAR(10)
    CHARACTER SET utf8mb4 COLLATE utf8mb4_unicode_ci NOT NULL,
  MODIFY col2 CHAR(10)
    CHARACTER SET utf8mb4 COLLATE utf8mb4_bin NOT NULL;
```

In terms of table content, conversion from `utf8` to `utf8mb4` presents no problems:

- For a BMP character, `utf8` and `utf8mb4` have identical storage characteristics: same code values, same encoding, same length.
- For a supplementary character, `utf8` cannot store the character at all, whereas `utf8mb4` requires four bytes. Because `utf8` cannot store the character at all, `utf8` columns have no supplementary characters and you need not worry about converting characters or losing data when converting to `utf8mb4`.

In terms of table structure, the catch when converting from `utf8` to `utf8mb4` is that the maximum length of a column or index key is unchanged in terms of *bytes*. Therefore, it is smaller in terms of *characters* because the maximum length of a character is four bytes instead of three. For the `CHAR`, `VARCHAR`, and `TEXT` data types, watch for these issues when converting your MySQL tables:

- Check all definitions of `utf8` columns and make sure they will not exceed the maximum length for the storage engine.
- Check all indexes on `utf8` columns and make sure they will not exceed the maximum length for the storage engine. Sometimes the maximum can change due to storage engine enhancements.

If the preceding conditions apply, you must either reduce the defined length of columns or indexes, or continue to use `utf8` rather than `utf8mb4`.

Here are some examples where structural changes may be needed:

- A `TINYTEXT` column can hold up to 255 bytes, so it can hold up to 85 3-byte or 63 4-byte characters. Suppose that you have a `TINYTEXT` column that uses `utf8` but must be able to contain more than 63 characters. You cannot convert it to `utf8mb4` unless you also change the data type to a longer type such as `TEXT`.

  Similarly, a very long `VARCHAR` column may need to be changed to one of the longer `TEXT` types if you want to convert it from `utf8` to `utf8mb4`.

- `InnoDB` has a maximum index length of 767 bytes for tables that use `COMPACT` or `REDUNDANT` row format, so for `utf8` or `utf8mb4` columns, you can index a maximum of 255 or 191 characters, respectively. If you currently have `utf8` columns with indexes longer than 191 characters, you must index a smaller number of characters.

  In an `InnoDB` table that uses `COMPACT` or `REDUNDANT` row format, these column and index definitions are legal:

  ```
  col1 VARCHAR(500) CHARACTER SET utf8, INDEX (col1(255))
  ```

  To use `utf8mb4` instead, the index must be smaller:

  ```
  col1 VARCHAR(500) CHARACTER SET utf8mb4, INDEX (col1(191))
  ```

Note

For `InnoDB` tables that use `COMPRESSED` or `DYNAMIC` row format, index key prefixes longer than 767 bytes (up to 3072 bytes) are permitted. Tables created with these row formats enable you to index a maximum of 1024 or 768 characters for `utf8` or `utf8mb4` columns, respectively. For related information, see Section 15.8.1.7, “Limits on InnoDB Tables”, and Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats”.

The preceding types of changes are most likely to be required only if you have very long columns or indexes. Otherwise, you should be able to convert your tables from `utf8` to `utf8mb4` without problems, using `ALTER TABLE` as described previously.

The following items summarize other potential areas of incompatibility:

- Performance of 4-byte UTF-8 (`utf8mb4`) is slower than for 3-byte UTF-8 (`utf8`). To avoid this penalty, continue to use `utf8`.
- `SET NAMES 'utf8mb4'` causes use of the 4-byte character set for connection character sets. As long as no 4-byte characters are sent from the server, there should be no problems. Otherwise, applications that expect to receive a maximum of three bytes per character may have problems. Conversely, applications that expect to send 4-byte characters must ensure that the server understands them. More generally, applications cannot send `utf8mb4`, `utf16`, `utf16le`, or `utf32` data to an older server that does not understand it:
  - `utf8mb4`, `utf16`, and `utf32` are not recognized before MySQL 5.5.3.
  - `utf16le` is not recognized before MySQL 5.6.1.
- For replication, if character sets that support supplementary characters are to be used on the master, all slaves must understand them as well. If you attempt to replicate from a newer master to an older slave, `utf8` data will be seen as `utf8` by the slave and should replicate correctly. But you cannot send `utf8mb4`, `utf16`, `utf16le`, or `utf32` data to an older slave that does not understand it:
  - `utf8mb4`, `utf16`, and `utf32` are not recognized before MySQL 5.5.3.
  - `utf16le` is not recognized before MySQL 5.6.1.

  Also, keep in mind the general principle that if a table has different definitions on the master and slave, this can lead to unexpected results. For example, the differences in maximum index key length make it risky to use `utf8` on the master and `utf8mb4` on the slave.

If you have converted to `utf8mb4`, `utf16`, `utf16le`, or `utf32`, and then decide to convert back to `utf8` or `ucs2` (for example, to downgrade to an older version of MySQL), these considerations apply:

- `utf8` and `ucs2` data should present no problems.
- The server must be recent enough to recognize definitions referring to the character set from which you are converting.
- For object definitions that refer to the `utf8mb4` character set, you can dump them with `mysqldump` prior to downgrading, edit the dump file to change instances of `utf8mb4` to `utf8`, and reload the file in the older server, as long as there are no 4-byte characters in the data. The older server will see `utf8` in the dump file object definitions and create new objects that use the (3-byte) `utf8` character set.

10.1.10 Supported Character Sets and Collations

This section indicates which character sets MySQL supports. There is one subsection for each group of related character sets. For each character set, the permissible collations are listed.

To list the available character sets and their default collations, use the `SHOW CHARACTER SET` statement or query the `INFORMATION_SCHEMA CHARACTER_SETS` table. For example:

```
mysql> SHOW CHARACTER SET;
+----------+---------------------------------+---------------------+--------+
| Charset  | Description                     | Default collation   | Maxlen |
+----------+---------------------------------+---------------------+--------+
| armscii8 | ARMSCII-8 Armenian              | armscii8_general_ci |      1 |
| ascii    | US ASCII                        | ascii_general_ci    |      1 |
| big5     | Big5 Traditional Chinese        | big5_chinese_ci     |      2 |
| binary   | Binary pseudo charset           | binary              |      1 |
| cp1250   | Windows Central European        | cp1250_general_ci   |      1 |
| cp1251   | Windows Cyrillic                | cp1251_general_ci   |      1 |
| cp1256   | Windows Arabic                  | cp1256_general_ci   |      1 |
| cp1257   | Windows Baltic                  | cp1257_general_ci   |      1 |
```

```
| cp850    | DOS West European                | cp850_general_ci    |      1 |
| cp852    | DOS Central European             | cp852_general_ci    |      1 |
| cp866    | DOS Russian                      | cp866_general_ci    |      1 |
| cp932    | SJIS for Windows Japanese        | cp932_japanese_ci   |      2 |
| dec8     | DEC West European                | dec8_swedish_ci     |      1 |
| eucjpms  | UJIS for Windows Japanese        | eucjpms_japanese_ci |      3 |
| euckr    | EUC-KR Korean                    | euckr_korean_ci     |      2 |
| gb18030  | China National Standard GB18030  | gb18030_chinese_ci  |      4 |
| gb2312   | GB2312 Simplified Chinese        | gb2312_chinese_ci   |      2 |
| gbk      | GBK Simplified Chinese           | gbk_chinese_ci      |      2 |
| geostd8  | GEOSTD8 Georgian                 | geostd8_general_ci  |      1 |
| greek    | ISO 8859-7 Greek                 | greek_general_ci    |      1 |
| hebrew   | ISO 8859-8 Hebrew                | hebrew_general_ci   |      1 |
| hp8      | HP West European                 | hp8_english_ci      |      1 |
| keybcs2  | DOS Kamenicky Czech-Slovak       | keybcs2_general_ci  |      1 |
| koi8r    | KOI8-R Relcom Russian            | koi8r_general_ci    |      1 |
| koi8u    | KOI8-U Ukrainian                 | koi8u_general_ci    |      1 |
| latin1   | cp1252 West European             | latin1_swedish_ci   |      1 |
| latin2   | ISO 8859-2 Central European      | latin2_general_ci   |      1 |
| latin5   | ISO 8859-9 Turkish               | latin5_turkish_ci   |      1 |
| latin7   | ISO 8859-13 Baltic               | latin7_general_ci   |      1 |
| macce    | Mac Central European             | macce_general_ci    |      1 |
| macroman | Mac West European                | macroman_general_ci |      1 |
| sjis     | Shift-JIS Japanese               | sjis_japanese_ci    |      2 |
| swe7     | 7bit Swedish                     | swe7_swedish_ci     |      1 |
| tis620   | TIS620 Thai                      | tis620_thai_ci      |      1 |
| ucs2     | UCS-2 Unicode                    | ucs2_general_ci     |      2 |
| ujis     | EUC-JP Japanese                  | ujis_japanese_ci    |      3 |
| utf16    | UTF-16 Unicode                   | utf16_general_ci    |      4 |
| utf16le  | UTF-16LE Unicode                 | utf16le_general_ci  |      4 |
| utf32    | UTF-32 Unicode                   | utf32_general_ci    |      4 |
| utf8     | UTF-8 Unicode                    | utf8_general_ci     |      3 |
| utf8mb4  | UTF-8 Unicode                    | utf8mb4_general_ci  |      4 |
+----------+----------------------------------+---------------------+--------+
```

In cases where a character set has multiple collations, it might not be clear which collation is most suitable for a given application. To avoid choosing the wrong collation, it can be helpful to perform some comparisons with representative data values to make sure that a given collation sorts values the way you expect.

10.1.10.1 Unicode Character Sets

MySQL supports multiple Unicode character sets:

- `utf8`, a UTF-8 encoding of the Unicode character set using one to three bytes per character.
- `ucs2`, the UCS-2 encoding of the Unicode character set using two bytes per character.
- `utf8mb4`, a UTF-8 encoding of the Unicode character set using one to four bytes per character.
- `utf16`, the UTF-16 encoding for the Unicode character set using two or four bytes per character. Like `ucs2` but with an extension for supplementary characters.
- `utf16le`, the UTF-16LE encoding for the Unicode character set. Like `utf16` but little-endian rather than big-endian.
- `utf32`, the UTF-32 encoding for the Unicode character set using four bytes per character.

`utf8` and `ucs2` support Basic Multilingual Plane (BMP) characters. `utf8mb4`, `utf16`, `utf16le`, and `utf32` support BMP and supplementary characters.

This section describes the collations available for Unicode character sets and their differentiating properties. For general information about Unicode, see Section 10.1.9, “Unicode Support”.

Most Unicode character sets have a general collation (indicated by `_general` in the name or by the absence of a language specifier), a binary collation (indicated by `_bin` in the name), and several language-specific collations (indicated by language specifiers). For example, for `utf8`,

`utf8_general_ci` and `utf8_bin` are its general and binary collations, and `utf8_danish_ci` is one of its language-specific collations.

Collation support for `utf16le` is limited. The only collations available are `utf16le_general_ci` and `utf16le_bin`. These are similar to `utf16_general_ci` and `utf16_bin`.

A locale code or language name shown in the following table indicates a language-specific collation. Unicode character sets may include collations for one or more of these languages.

Table 10.3 Unicode Collation Language Specifiers

Language	Language Specifier
Classical Latin	`la` or `roman`
Croatian	`hr` or `croatian`
Czech	`cs` or `czech`
Danish	`da` or `danish`
Esperanto	`eo` or `esperanto`
Estonian	`et` or `estonian`
German phone book order	`de_pb` or `german2`
Hungarian	`hu` or `hungarian`
Icelandic	`is` or `icelandic`
Japanese	`ja`
Latvian	`lv` or `latvian`
Lithuanian	`lt` or `lithuanian`
Persian	`persian`
Polish	`pl` or `polish`
Romanian	`ro` or `romanian`
Sinhala	`sinhala`
Slovak	`sk` or `slovak`
Slovenian	`sl` or `slovenian`
Modern Spanish	`es` or `spanish`
Traditional Spanish	`es_trad` or `spanish2`
Swedish	`sv` or `swedish`
Turkish	`tr` or `turkish`
Vietnamese	`vi` or `vietnamese`

Croatian collations are tailored for these Croatian letters: `Č`, `Ć`, `Dž`, `Đ`, `Lj`, `Nj`, `Š`, `Ž`.

Danish collations may also be used for Norwegian.

For Japanese, the `utf8mb4` character set includes `utf8mb4_ja_0900_as_cs` and `utf8mb4_ja_0900_as_cs_ks` collations. Both collations are accent sensitive and case sensitive. `utf8mb4_ja_0900_as_cs_ks` is also kana sensitive and distinguishes Katakana characters from Hiragana characters, whereas `utf8mb4_ja_0900_as_cs` treats Katakana and Hiragana characters as equal for sorting. Applications that require a Japanese collation but not kana sensitivity may use `utf8mb4_ja_0900_as_cs` for better sort performance. `utf8mb4_ja_0900_as_cs` uses three weight levels for sorting; `utf8mb4_ja_0900_as_cs_ks` uses four.

For Classical Latin collations that are accent insensitive, `I` and `J` compare as equal, and `U` and `V` compare as equal. `I` and `J`, and `U` and `V` compare as equal on the base letter level. In other words, `J` is regarded as an accented `I`, and `U` is regarded as an accented `V`.

Spanish collations are available for modern and traditional Spanish. For both, `ñ` (n-tilde) is a separate letter between `n` and `o`. In addition, for traditional Spanish, `ch` is a separate letter between `c` and `d`, and `ll` is a separate letter between `l` and `m`.

Traditional Spanish collations may also be used for Asturian and Galician.

Swedish collations include Swedish rules. For example, in Swedish, the following relationship holds, which is not something expected by a German or French speaker:

```
Ü = Y < Ö
```

For questions about particular language orderings, unicode.org provides Common Locale Data Repository (CLDR) collation charts at http://www.unicode.org/cldr/charts/30/collation/index.html.

The *`xxx`*`_general_mysql500_ci` collations preserve the pre-5.1.24 ordering of the original *`xxx`*`_general_ci` collations and permit upgrades for tables created before MySQL 5.1.24 (Bug #27877).

MySQL implements the *`xxx`*`_unicode_ci` collations according to the Unicode Collation Algorithm (UCA) described at http://www.unicode.org/reports/tr10/. The collation uses the version-4.0.0 UCA weight keys: http://www.unicode.org/Public/UCA/4.0.0/allkeys-4.0.0.txt. The *`xxx`*`_unicode_ci` collations have only partial support for the Unicode Collation Algorithm. Some characters are not supported, and combining marks are not fully supported. This affects primarily Vietnamese, Yoruba, and some smaller languages such as Navajo. A combined character is considered different from the same character written with a single unicode character in string comparisons, and the two characters are considered to have a different length (for example, as returned by the `CHAR_LENGTH()` function or in result set metadata).

Unicode collations based on UCA versions later than 4.0.0 include the version in the collation name. Thus, `utf8mb4_unicode_520_ci` is based on UCA 5.2.0 weight keys (http://www.unicode.org/Public/UCA/5.2.0/allkeys.txt), whereas `utf8mb4_0900_ai_ci` is based on UCA 9.0.0 weight keys (http://www.unicode.org/Public/UCA/9.0.0/allkeys.txt). For collations of `utf8` that include a UCA version, there is no `utf8mb3` alias; see Section 10.1.9.2, “The utf8mb3 Character Set (Alias for utf8)”.

Collations based on UCA 9.0.0 and higher are faster than collations based on UCA versions below 9.0.0. They also have a pad attribute of NO PAD, in contrast to PAD SPACE as used in collations based on UCA versions below 9.0.0. NO PAD collations treat spaces at the end of strings like any other character.

To determine the pad attribute for a collation, use the `INFORMATION_SCHEMA COLLATIONS` table, which has a `PAD_ATTRIBUTE` column.

Comparisons of `VARCHAR` columns that have a NO PAD collation differ from other collations with respect to trailing spaces. For example, `'a'` and `'a '` compare as different strings, not the same string.

MySQL implements language-specific Unicode collations if the ordering based only on UCA does not work well for a language. Language-specific collations are UCA-based, with additional language tailoring rules.

For example, the nonlanguage-specific `utf8mb4_0900_ai_ci` and language-specific `utf8mb4_`*`LOCALE`*`_0900_ai_ci` Unicode collations each have these characteristics:

- The collation is based on Unicode Collation Algorithm (UCA) 9.0.0 and Common Locale Data Repository (CLDR) v30, is accent insensitive, and case insensitive. These characteristics are indicated by `_0900`, `_ai`, and `_ci` in the collation name. Exception: `utf8mb4_la_0900_ai_ci` is not based on CLDR because Classical Latin is not defined in CLDR.

- The collation works for all characters in the range [U+0, U+10FFFF].

- If the collation is not language specific, it sorts all characters, including supplementary characters, in default order (described following). If the collation is language specific, it sorts characters of the language correctly according to language-specific rules, and characters not in the language in default order.
- By default, the collation sorts characters having a code point listed in the DUCET table (Default Unicode Collation Element Table) according to the weight value assigned in the table. The collation sorts characters not having a code point listed in the DUCET table using their implicit weight value, which is constructed according to the UCA.
- For non-language-specific collations, characters in contraction sequences are treated as separate characters. For language-specific collations, contractions might change character sorting order.

`LOWER()` and `UPPER()` perform case folding according to the collation of their argument. A character that has uppercase and lowercase versions only in a Unicode version more recent than 4.0.0 is converted by these functions only if the argument has a collation that uses a recent enough UCA version.

For any Unicode character set, operations performed using the *`xxx`*`_general_ci` collation are faster than those for the *`xxx`*`_unicode_ci` collation. For example, comparisons for the `utf8_general_ci` collation are faster, but slightly less correct, than comparisons for `utf8_unicode_ci`. The reason for this is that `utf8_unicode_ci` supports mappings such as expansions; that is, when one character compares as equal to combinations of other characters. For example, in German and some other languages ß is equal to `ss`. `utf8_unicode_ci` also supports contractions and ignorable characters. `utf8_general_ci` is a legacy collation that does not support expansions, contractions, or ignorable characters. It can make only one-to-one comparisons between characters.

To further illustrate, the following equalities hold in both `utf8_general_ci` and `utf8_unicode_ci` (for the effect of this in comparisons or searches, see Section 10.1.8.6, “Examples of the Effect of Collation”):

```
Ä = A
Ö = O
Ü = U
```

A difference between the collations is that this is true for `utf8_general_ci`:

```
ß = s
```

Whereas this is true for `utf8_unicode_ci`, which supports the German DIN-1 ordering (also known as dictionary order):

```
ß = ss
```

MySQL implements `utf8` language-specific collations if the ordering with `utf8_unicode_ci` does not work well for a language. For example, `utf8_unicode_ci` works fine for German dictionary order and French, so there is no need to create special `utf8` collations.

`utf8_general_ci` also is satisfactory for both German and French, except that ß is equal to `s`, and not to `ss`. If this is acceptable for your application, you should use `utf8_general_ci` because it is faster. If this is not acceptable (for example, if you require German dictionary order), use `utf8_unicode_ci` because it is more accurate.

If you require German DIN-2 (phone book) ordering, use the `utf8_german2_ci` collation, which compares the following sets of characters equal:

```
Ä = Æ = AE
Ö = Œ = OE
```

```
Ü = UE
ß = ss
```

`utf8_german2_ci` is similar to `latin1_german2_ci`, but the latter does not compare `Æ` equal to `AE` or `Œ` equal to `OE`. There is no `utf8_german_ci` corresponding to `latin1_german_ci` for German dictionary order because `utf8_general_ci` suffices.

For all Unicode collations except the binary (`_bin`) collations, MySQL performs a table lookup to find a character's collating weight. This weight can be displayed using the `WEIGHT_STRING()` function. (See Section 12.5, “String Functions”.) If a character is not in the table (for example, because it is a “new” character), collating weight determination becomes more complex:

- For BMP characters in general collations (*`xxx`*`_general_ci`), weight = code point.
- For BMP characters in UCA collations (for example, *`xxx`*`_unicode_ci` and language-specific collations), the following algorithm applies:

  ```
  if (code >= 0x3400 && code <= 0x4DB5)
    base= 0xFB80; /* CJK Ideograph Extension */
  else if (code >= 0x4E00 && code <= 0x9FA5)
    base= 0xFB40; /* CJK Ideograph */
  else
    base= 0xFBC0; /* All other characters */
  aaaa= base +  (code >> 15);
  bbbb= (code & 0x7FFF) | 0x8000;
  ```

  The result is a sequence of two collating elements, `aaaa` followed by `bbbb`. For example:

  ```
  mysql> SELECT HEX(WEIGHT_STRING(_ucs2 0x04CF COLLATE ucs2_unicode_ci));
  +-------------------------------------------------------------+
  | HEX(WEIGHT_STRING(_ucs2 0x04CF COLLATE ucs2_unicode_ci)) |
  +-------------------------------------------------------------+
  | FBC084CF                                                    |
  +-------------------------------------------------------------+
  ```

  Thus, `U+04cf CYRILLIC SMALL LETTER PALOCHKA` is, with all UCA 4.0.0 collations, greater than `U+04c0 CYRILLIC LETTER PALOCHKA`. With UCA 5.2.0 collations, all palochkas sort together.

- For supplementary characters in general collations, the weight is the weight for `0xfffd REPLACEMENT CHARACTER`. For supplementary characters in UCA 4.0.0 collations, their collating weight is `0xfffd`. That is, to MySQL, all supplementary characters are equal to each other, and greater than almost all BMP characters.

  An example with Deseret characters and `COUNT(DISTINCT)`:

  ```
  CREATE TABLE t (s1 VARCHAR(5) CHARACTER SET utf32 COLLATE utf32_unicode_ci);
  INSERT INTO t VALUES (0xfffd);   /* REPLACEMENT CHARACTER */
  INSERT INTO t VALUES (0x010412); /* DESERET CAPITAL LETTER BEE */
  INSERT INTO t VALUES (0x010413); /* DESERET CAPITAL LETTER TEE */
  SELECT COUNT(DISTINCT s1) FROM t;
  ```

  The result is 2 because in the MySQL *`xxx`*`_unicode_ci` collations, the replacement character has a weight of `0x0dc6`, whereas Deseret Bee and Deseret Tee both have a weight of `0xfffd`. (Were the `utf32_general_ci` collation used instead, the result is 1 because all three characters have a weight of `0xfffd` in that collation.)

  An example with cuneiform characters and `WEIGHT_STRING()`:

  ```
  /*
  The four characters in the INSERT string are
  00000041  # LATIN CAPITAL LETTER A
  ```

```
0001218F  # CUNEIFORM SIGN KAB
000121A7  # CUNEIFORM SIGN KISH
00000042  # LATIN CAPITAL LETTER B
*/
CREATE TABLE t (s1 CHAR(4) CHARACTER SET utf32 COLLATE utf32_unicode_ci);
INSERT INTO t VALUES (0x000004100001218f000121a700000042);
SELECT HEX(WEIGHT_STRING(s1)) FROM t;
```

The result is:

```
0E33 FFFD FFFD 0E4A
```

`0E33` and `0E4A` are primary weights as in UCA 4.0.0. `FFFD` is the weight for KAB and also for KISH.

The rule that all supplementary characters are equal to each other is nonoptimal but is not expected to cause trouble. These characters are very rare, so it is very rare that a multi-character string consists entirely of supplementary characters. In Japan, since the supplementary characters are obscure Kanji ideographs, the typical user does not care what order they are in, anyway. If you really want rows sorted by the MySQL rule and secondarily by code point value, it is easy:

```
ORDER BY s1 COLLATE utf32_unicode_ci, s1 COLLATE utf32_bin
```

- For supplementary characters based on UCA versions higher than 4.0.0 (for example, *`xxx`*`_unicode_520_ci`), supplementary characters do not necessarily all have the same collation weight. Some have explicit weights from the UCA `allkeys.txt` file. Others have weights calculated from this algorithm:

```
aaaa= base +  (code >> 15);
bbbb= (code & 0x7FFF) | 0x8000;
```

There is a difference between “ordering by the character's code value” and “ordering by the character's binary representation,” a difference that appears only with `utf16_bin`, because of surrogates.

Suppose that `utf16_bin` (the binary collation for `utf16`) was a binary comparison “byte by byte” rather than “character by character.” If that were so, the order of characters in `utf16_bin` would differ from the order in `utf8_bin`. For example, the following chart shows two rare characters. The first character is in the range `E000`-`FFFF`, so it is greater than a surrogate but less than a supplementary. The second character is a supplementary.

```
Code point  Character                   utf8         utf16
----------  ---------                   ----         -----
0FF9D       HALFWIDTH KATAKANA LETTER N  EF BE 9D     FF 9D
10384       UGARITIC LETTER DELTA        F0 90 8E 84  D8 00 DF 84
```

The two characters in the chart are in order by code point value because `0xff9d` < `0x10384`. And they are in order by `utf8` value because `0xef` < `0xf0`. But they are not in order by `utf16` value, if we use byte-by-byte comparison, because `0xff` > `0xd8`.

So MySQL's `utf16_bin` collation is not “byte by byte.” It is “by code point.” When MySQL sees a supplementary-character encoding in `utf16`, it converts to the character's code-point value, and then compares. Therefore, `utf8_bin` and `utf16_bin` are the same ordering. This is consistent with the SQL:2008 standard requirement for a UCS_BASIC collation: “UCS_BASIC is a collation in which the ordering is determined entirely by the Unicode scalar values of the characters in the strings being sorted. It is applicable to the UCS character repertoire. Since every character repertoire is a subset of the UCS repertoire, the UCS_BASIC collation is potentially applicable to every character set. NOTE 11: The Unicode scalar value of a character is its code point treated as an unsigned integer.”

If the character set is `ucs2`, comparison is byte-by-byte, but `ucs2` strings should not contain surrogates, anyway.

10.1.10.2 West European Character Sets

Western European character sets cover most West European languages, such as French, Spanish, Catalan, Basque, Portuguese, Italian, Albanian, Dutch, German, Danish, Swedish, Norwegian, Finnish, Faroese, Icelandic, Irish, Scottish, and English.

- `ascii` (US ASCII) collations:
  - `ascii_bin`
  - `ascii_general_ci` (default)
- `cp850` (DOS West European) collations:
  - `cp850_bin`
  - `cp850_general_ci` (default)
- `dec8` (DEC Western European) collations:
  - `dec8_bin`
  - `dec8_swedish_ci` (default)
- `hp8` (HP Western European) collations:
  - `hp8_bin`
  - `hp8_english_ci` (default)
- `latin1` (cp1252 West European) collations:
  - `latin1_bin`
  - `latin1_danish_ci`
  - `latin1_general_ci`
  - `latin1_general_cs`
  - `latin1_german1_ci`
  - `latin1_german2_ci`
  - `latin1_spanish_ci`
  - `latin1_swedish_ci` (default)

  `latin1` is the default character set. MySQL's `latin1` is the same as the Windows `cp1252` character set. This means it is the same as the official `ISO 8859-1` or IANA (Internet Assigned Numbers Authority) `latin1`, except that IANA `latin1` treats the code points between `0x80` and `0x9f` as "undefined," whereas `cp1252`, and therefore MySQL's `latin1`, assign characters for those positions. For example, `0x80` is the Euro sign. For the "undefined" entries in `cp1252`, MySQL translates `0x81` to Unicode `0x0081`, `0x8d` to `0x008d`, `0x8f` to `0x008f`, `0x90` to `0x0090`, and `0x9d` to `0x009d`.

  The `latin1_swedish_ci` collation is the default that probably is used by the majority of MySQL customers. Although it is frequently said that it is based on the Swedish/Finnish collation rules, there are Swedes and Finns who disagree with this statement.

  The `latin1_german1_ci` and `latin1_german2_ci` collations are based on the DIN-1 and DIN-2 standards, where DIN stands for *Deutsches Institut für Normung* (the German equivalent of ANSI). DIN-1 is called the "dictionary collation" and DIN-2 is called the "phone book collation." For

an example of the effect this has in comparisons or when doing searches, see Section 10.1.8.6, “Examples of the Effect of Collation”.

- `latin1_german1_ci` (dictionary) rules:

```
Ä = A
Ö = O
Ü = U
ß = s
```

- `latin1_german2_ci` (phone-book) rules:

```
Ä = AE
Ö = OE
Ü = UE
ß = ss
```

In the `latin1_spanish_ci` collation, `ñ` (n-tilde) is a separate letter between `n` and `o`.

- `macroman` (Mac West European) collations:
  - `macroman_bin`
  - `macroman_general_ci` (default)
- `swe7` (7bit Swedish) collations:
  - `swe7_bin`
  - `swe7_swedish_ci` (default)

10.1.10.3 Central European Character Sets

MySQL provides some support for character sets used in the Czech Republic, Slovakia, Hungary, Romania, Slovenia, Croatia, Poland, and Serbia (Latin).

- `cp1250` (Windows Central European) collations:
  - `cp1250_bin`
  - `cp1250_croatian_ci`
  - `cp1250_czech_cs`
  - `cp1250_general_ci` (default)
  - `cp1250_polish_ci`
- `cp852` (DOS Central European) collations:
  - `cp852_bin`
  - `cp852_general_ci` (default)
- `keybcs2` (DOS Kamenicky Czech-Slovak) collations:
  - `keybcs2_bin`
  - `keybcs2_general_ci` (default)
- `latin2` (ISO 8859-2 Central European) collations:
  - `latin2_bin`

  - `latin2_croatian_ci`
  - `latin2_czech_cs`
  - `latin2_general_ci` (default)
  - `latin2_hungarian_ci`
- `macce` (Mac Central European) collations:
  - `macce_bin`
  - `macce_general_ci` (default)

10.1.10.4 South European and Middle East Character Sets

South European and Middle Eastern character sets supported by MySQL include Armenian, Arabic, Georgian, Greek, Hebrew, and Turkish.

- `armscii8` (ARMSCII-8 Armenian) collations:
  - `armscii8_bin`
  - `armscii8_general_ci` (default)
- `cp1256` (Windows Arabic) collations:
  - `cp1256_bin`
  - `cp1256_general_ci` (default)
- `geostd8` (GEOSTD8 Georgian) collations:
  - `geostd8_bin`
  - `geostd8_general_ci` (default)
- `greek` (ISO 8859-7 Greek) collations:
  - `greek_bin`
  - `greek_general_ci` (default)
- `hebrew` (ISO 8859-8 Hebrew) collations:
  - `hebrew_bin`
  - `hebrew_general_ci` (default)
- `latin5` (ISO 8859-9 Turkish) collations:
  - `latin5_bin`
  - `latin5_turkish_ci` (default)

10.1.10.5 Baltic Character Sets

The Baltic character sets cover Estonian, Latvian, and Lithuanian languages.

- `cp1257` (Windows Baltic) collations:
  - `cp1257_bin`

  - `cp1257_general_ci` (default)
  - `cp1257_lithuanian_ci`
- `latin7` (ISO 8859-13 Baltic) collations:
  - `latin7_bin`
  - `latin7_estonian_cs`
  - `latin7_general_ci` (default)
  - `latin7_general_cs`

10.1.10.6 Cyrillic Character Sets

The Cyrillic character sets and collations are for use with Belarusian, Bulgarian, Russian, Ukrainian, and Serbian (Cyrillic) languages.

- `cp1251` (Windows Cyrillic) collations:
  - `cp1251_bin`
  - `cp1251_bulgarian_ci`
  - `cp1251_general_ci` (default)
  - `cp1251_general_cs`
  - `cp1251_ukrainian_ci`
- `cp866` (DOS Russian) collations:
  - `cp866_bin`
  - `cp866_general_ci` (default)
- `koi8r` (KOI8-R Relcom Russian) collations:
  - `koi8r_bin`
  - `koi8r_general_ci` (default)
- `koi8u` (KOI8-U Ukrainian) collations:
  - `koi8u_bin`
  - `koi8u_general_ci` (default)

10.1.10.7 Asian Character Sets

The Asian character sets that we support include Chinese, Japanese, Korean, and Thai. These can be complicated. For example, the Chinese sets must allow for thousands of different characters. See The cp932 Character Set, for additional information about the `cp932` and `sjis` character sets. See The gb18030 Character Set, for additional information about character set support for the Chinese National Standard GB 18030.

For answers to some common questions and problems relating support for Asian character sets in MySQL, see Section A.11, “MySQL 8.0 FAQ: MySQL Chinese, Japanese, and Korean Character Sets”.

- `big5` (Big5 Traditional Chinese) collations:
  - `big5_bin`

    - `big5_chinese_ci` (default)
- `cp932` (SJIS for Windows Japanese) collations:
    - `cp932_bin`
    - `cp932_japanese_ci` (default)
- `eucjpms` (UJIS for Windows Japanese) collations:
    - `eucjpms_bin`
    - `eucjpms_japanese_ci` (default)
- `euckr` (EUC-KR Korean) collations:
    - `euckr_bin`
    - `euckr_korean_ci` (default)
- `gb2312` (GB2312 Simplified Chinese) collations:
    - `gb2312_bin`
    - `gb2312_chinese_ci` (default)
- `gbk` (GBK Simplified Chinese) collations:
    - `gbk_bin`
    - `gbk_chinese_ci` (default)
- `gb18030` (China National Standard GB18030) collations:
    - `gb18030_bin`
    - `gb18030_chinese_ci` (default)
    - `gb18030_unicode_520_ci`
- `sjis` (Shift-JIS Japanese) collations:
    - `sjis_bin`
    - `sjis_japanese_ci` (default)
- `tis620` (TIS620 Thai) collations:
    - `tis620_bin`
    - `tis620_thai_ci` (default)
- `ujis` (EUC-JP Japanese) collations:
    - `ujis_bin`
    - `ujis_japanese_ci` (default)

The `big5_chinese_ci` collation sorts on number of strokes.

The cp932 Character Set

Why is `cp932` needed?

In MySQL, the `sjis` character set corresponds to the `Shift_JIS` character set defined by IANA, which supports JIS X0201 and JIS X0208 characters. (See http://www.iana.org/assignments/character-sets.)

However, the meaning of “SHIFT JIS” as a descriptive term has become very vague and it often includes the extensions to `Shift_JIS` that are defined by various vendors.

For example, “SHIFT JIS” used in Japanese Windows environments is a Microsoft extension of `Shift_JIS` and its exact name is `Microsoft Windows Codepage : 932` or `cp932`. In addition to the characters supported by `Shift_JIS`, `cp932` supports extension characters such as NEC special characters, NEC selected—IBM extended characters, and IBM selected characters.

Many Japanese users have experienced problems using these extension characters. These problems stem from the following factors:

- MySQL automatically converts character sets.
- Character sets are converted using Unicode (`ucs2`).
- The `sjis` character set does not support the conversion of these extension characters.
- There are several conversion rules from so-called “SHIFT JIS” to Unicode, and some characters are converted to Unicode differently depending on the conversion rule. MySQL supports only one of these rules (described later).

The MySQL `cp932` character set is designed to solve these problems.

Because MySQL supports character set conversion, it is important to separate IANA `Shift_JIS` and `cp932` into two different character sets because they provide different conversion rules.

How does `cp932` differ from `sjis`?

The `cp932` character set differs from `sjis` in the following ways:

- `cp932` supports NEC special characters, NEC selected—IBM extended characters, and IBM selected characters.
- Some `cp932` characters have two different code points, both of which convert to the same Unicode code point. When converting from Unicode back to `cp932`, one of the code points must be selected. For this “round trip conversion,” the rule recommended by Microsoft is used. (See http://support.microsoft.com/kb/170559/EN-US/.)

  The conversion rule works like this:

  - If the character is in both JIS X 0208 and NEC special characters, use the code point of JIS X 0208.
  - If the character is in both NEC special characters and IBM selected characters, use the code point of NEC special characters.
  - If the character is in both IBM selected characters and NEC selected—IBM extended characters, use the code point of IBM extended characters.

  The table shown at https://msdn.microsoft.com/en-us/goglobal/cc305152.aspx provides information about the Unicode values of `cp932` characters. For `cp932` table entries with characters under which a four-digit number appears, the number represents the corresponding Unicode (`ucs2`) encoding. For table entries with an underlined two-digit value appears, there is a range of `cp932` character values that begin with those two digits. Clicking such a table entry takes you to a page that displays the Unicode value for each of the `cp932` characters that begin with those digits.

  The following links are of special interest. They correspond to the encodings for the following sets of characters:

- NEC special characters (lead byte `0x87`):

```
https://msdn.microsoft.com/en-us/goglobal/gg674964
```

- NEC selected—IBM extended characters (lead byte `0xED` and `0xEE`):

```
https://msdn.microsoft.com/en-us/goglobal/gg671837
https://msdn.microsoft.com/en-us/goglobal/gg671838
```

- IBM selected characters (lead byte `0xFA`, `0xFB`, `0xFC`):

```
https://msdn.microsoft.com/en-us/goglobal/gg671839
https://msdn.microsoft.com/en-us/goglobal/gg671840
https://msdn.microsoft.com/en-us/goglobal/gg671841
```

- `cp932` supports conversion of user-defined characters in combination with `eucjpms`, and solves the problems with `sjis`/`ujis` conversion. For details, please refer to http://www.sljfaq.org/afaq/encodings.html.

For some characters, conversion to and from `ucs2` is different for `sjis` and `cp932`. The following tables illustrate these differences.

Conversion to `ucs2`:

`sjis`/`cp932` Value	**`sjis` -> `ucs2` Conversion**	**`cp932` -> `ucs2` Conversion**
5C	005C	005C
7E	007E	007E
815C	2015	2015
815F	005C	FF3C
8160	301C	FF5E
8161	2016	2225
817C	2212	FF0D
8191	00A2	FFE0
8192	00A3	FFE1
81CA	00AC	FFE2

Conversion from `ucs2`:

`ucs2` value	**`ucs2` -> `sjis` Conversion**	**`ucs2` -> `cp932` Conversion**
005C	815F	5C
007E	7E	7E
00A2	8191	3F
00A3	8192	3F
00AC	81CA	3F
2015	815C	815C
2016	8161	3F
2212	817C	3F
2225	3F	8161
301C	8160	3F
FF0D	3F	817C

ucs2 value	ucs2 -> sjis Conversion	ucs2 -> cp932 Conversion
FF3C	3F	815F
FF5E	3F	8160
FFE0	3F	8191
FFE1	3F	8192
FFE2	3F	81CA

Users of any Japanese character sets should be aware that using `--character-set-client-handshake` (or `--skip-character-set-client-handshake`) has an important effect. See Section 5.1.4, “Server Command Options”.

The gb18030 Character Set

In MySQL, the `gb18030` character set corresponds to the “Chinese National Standard GB 18030-2005: Information technology—Chinese coded character set”, which is the official character set of the People's Republic of China (PRC).

Characteristics of the MySQL gb18030 Character Set

- Supports all code points defined by the GB 18030-2005 standard. Unassigned code points in the ranges (GB+8431A439, GB+90308130) and (GB+E3329A36, GB+EF39EF39) are treated as '?' (0x3F). Conversion of unassigned code points return '?'.
- Supports UPPER and LOWER conversion for all GB18030 code points. Case folding defined by Unicode is also supported (based on `CaseFolding-6.3.0.txt`).
- Supports Conversion of data to and from other character sets.
- Supports SQL statements such as `SET NAMES`.
- Supports comparison between `gb18030` strings, and between `gb18030` strings and strings of other character sets. There is a conversion if strings have different character sets. Comparisons that include or ignore trailing spaces are also supported.
- The private use area (U+E000, U+F8FF) in Unicode is mapped to `gb18030`.
- There is no mapping between (U+D800, U+DFFF) and GB18030. Attempted conversion of code points in this range returns '?'.
- If an incoming sequence is illegal, an error or warning is returned. If an illegal sequence is used in `CONVERT()`, an error is returned. Otherwise, a warning is returned.
- For consistency with `utf8` and `utf8mb4`, UPPER is not supported for ligatures.
- Searches for ligatures also match uppercase ligatures when using the `gb18030_unicode_520_ci` collation.
- If a character has more than one uppercase character, the chosen uppercase character is the one whose lowercase is the character itself.
- The minimum multibyte length is 1 and the maximum is 4. The character set determines the length of a sequence using the first 1 or 2 bytes.

Supported Collations

- `gb18030_bin`: A binary collation.
- `gb18030_chinese_ci`: The default collation, which supports Pinyin. Sorting of non-Chinese characters is based on the order of the original sort key. The original sort key is `GB(UPPER(ch))`

if `UPPER(ch)` exists. Otherwise, the original sort key is `GB(ch)`. Chinese characters are sorted according to the Pinyin collation defined in the Unicode Common Locale Data Repository (CLDR 24). Non-Chinese characters are sorted before Chinese characters with the exception of `GB+FE39FE39`, which is the code point maximum.

- `gb18030_unicode_520_ci`: A Unicode collation. Use this collation if you need to ensure that ligatures are sorted correctly.

10.1.10.8 The Binary Character Set

The `binary` character set is the chararcter set of binary strings, which are sequences of bytes. The `binary` character set has one collation, also named `binary`. Comparison and sorting are based on numeric byte values. The effect is that lettercase and accent differences are significant in comparisons. That is, the `binary` collation is case sensitive and accent sensitive.

```
mysql> SET NAMES 'binary';
mysql> SELECT CHARSET('abc'), COLLATION('abc');
+----------------+------------------+
| CHARSET('abc') | COLLATION('abc') |
+----------------+------------------+
| binary         | binary           |
+----------------+------------------+
mysql> SELECT 'abc' = 'ABC', 'a' = 'ä';
+--------------+------------+
| 'abc' = 'ABC' | 'a' = 'ä'  |
+--------------+------------+
|             0 |          0 |
+--------------+------------+
```

For information about the differences between the `binary` collation of the `binary` character set and the `_bin` collations of nonbinary character sets, see Section 10.1.8.5, “The binary Collation Compared to _bin Collations”.

To convert a string expression to a binary string, any of these constructs are equivalent:

```
BINARY expr
CAST(expr AS BINARY)
CONVERT(expr USING BINARY)
```

If *`expr`* is a character string literal, the `_binary` introducer may be used to designate it as a binary string. For example:

```
_binary 'a'
```

The `_binary` introducer is permitted for hexadecimal literals and bit-value literals as well, but unnecessary; such literals are binary strings by default.

For more information about introducers, see Section 10.1.3.8, “Character Set Introducers”.

10.2 Setting the Error Message Language

By default, `mysqld` produces error messages in English, but they can also be displayed in any of several other languages: Czech, Danish, Dutch, Estonian, French, German, Greek, Hungarian, Italian, Japanese, Korean, Norwegian, Norwegian-ny, Polish, Portuguese, Romanian, Russian, Slovak, Spanish, or Swedish.

You can select which language the server uses for error messages using the instructions in this section.

The server searches for the error message file in two locations:

- It tries to find the file in a directory constructed from two system variable values, `lc_messages_dir` and `lc_messages`, with the latter converted to a language name. Suppose that you start the server using this command:

```
shell> mysqld --lc_messages_dir=/usr/share/mysql --lc_messages=fr_FR
```

  In this case, `mysqld` maps the locale `fr_FR` to the language `french` and looks for the error file in the `/usr/share/mysql/french` directory.

- If the message file cannot be found in the directory constructed as just described, the server ignores the `lc_messages` value and uses only the `lc_messages_dir` value as the location in which to look.

The `lc_messages_dir` system variable has only a global value and is read only. `lc_messages` has global and session values and can be modified at runtime, so the error message language can be changed while the server is running, and individual clients each can have a different error message language by changing their session `lc_messages` value to a different locale name. For example, if the server is using the `fr_FR` locale for error messages, a client can execute this statement to receive error messages in English:

```
mysql> SET lc_messages = 'en_US';
```

By default, the language files are located in the `share/mysql/LANGUAGE` directory under the MySQL base directory.

For information about changing the character set for error messages (rather than the language), see Section 10.1.6, “Error Message Character Set”.

You can change the content of the error messages produced by the server using the instructions in the MySQL Internals manual, available at MySQL Internals: Error Messages. If you do change the content of error messages, remember to repeat your changes after each upgrade to a newer version of MySQL.

10.3 Adding a Character Set

This section discusses the procedure for adding a character set to MySQL. The proper procedure depends on whether the character set is simple or complex:

- If the character set does not need special string collating routines for sorting and does not need multibyte character support, it is simple.
- If the character set needs either of those features, it is complex.

For example, `greek` and `swe7` are simple character sets, whereas `big5` and `czech` are complex character sets.

To use the following instructions, you must have a MySQL source distribution. In the instructions, *`MYSET`* represents the name of the character set that you want to add.

1. Add a `<charset>` element for *`MYSET`* to the `sql/share/charsets/Index.xml` file. Use the existing contents in the file as a guide to adding new contents. A partial listing for the `latin1` `<charset>` element follows:

```
<charset name="latin1">
  <family>Western</family>
  <description>cp1252 West European</description>
  ...
  <collation name="latin1_swedish_ci" id="8" order="Finnish, Swedish">
    <flag>primary</flag>
    <flag>compiled</flag>
```

```
  </collation>
  <collation name="latin1_danish_ci" id="15" order="Danish"/>
  ...
  <collation name="latin1_bin" id="47" order="Binary">
    <flag>binary</flag>
    <flag>compiled</flag>
  </collation>
  ...
</charset>
```

The `<charset>` element must list all the collations for the character set. These must include at least a binary collation and a default (primary) collation. The default collation is often named using a suffix of `general_ci` (general, case insensitive). It is possible for the binary collation to be the default collation, but usually they are different. The default collation should have a `primary` flag. The binary collation should have a `binary` flag.

You must assign a unique ID number to each collation. The range of IDs from 1024 to 2047 is reserved for user-defined collations. To find the maximum of the currently used collation IDs, use this query:

```
SELECT MAX(ID) FROM INFORMATION_SCHEMA.COLLATIONS;
```

2. This step depends on whether you are adding a simple or complex character set. A simple character set requires only a configuration file, whereas a complex character set requires C source file that defines collation functions, multibyte functions, or both.

   For a simple character set, create a configuration file, *`MYSET`*`.xml`, that describes the character set properties. Create this file in the `sql/share/charsets` directory. You can use a copy of `latin1.xml` as the basis for this file. The syntax for the file is very simple:

   - Comments are written as ordinary XML comments (`<!-- text -->`).
   - Words within `<map>` array elements are separated by arbitrary amounts of whitespace.
   - Each word within `<map>` array elements must be a number in hexadecimal format.
   - The `<map>` array element for the `<ctype>` element has 257 words. The other `<map>` array elements after that have 256 words. See Section 10.3.1, “Character Definition Arrays”.
   - For each collation listed in the `<charset>` element for the character set in `Index.xml`, *`MYSET`*`.xml` must contain a `<collation>` element that defines the character ordering.

   For a complex character set, create a C source file that describes the character set properties and defines the support routines necessary to properly perform operations on the character set:

   - Create the file `ctype-`*`MYSET`*`.c` in the `strings` directory. Look at one of the existing `ctype-*.c` files (such as `ctype-big5.c`) to see what needs to be defined. The arrays in your file must have names like `ctype_`*`MYSET`*, `to_lower_`*`MYSET`*, and so on. These correspond to the arrays for a simple character set. See Section 10.3.1, “Character Definition Arrays”.
   - For each `<collation>` element listed in the `<charset>` element for the character set in `Index.xml`, the `ctype-`*`MYSET`*`.c` file must provide an implementation of the collation.
   - If the character set requires string collating functions, see Section 10.3.2, “String Collating Support for Complex Character Sets”.
   - If the character set requires multibyte character support, see Section 10.3.3, “Multi-Byte Character Support for Complex Character Sets”.

3. Modify the configuration information. Use the existing configuration information as a guide to adding information for *`MYSYS`*. The example here assumes that the character set has default and binary collations, but more lines are needed if *`MYSET`* has additional collations.

a. Edit `mysys/charset-def.c`, and “register” the collations for the new character set.

   Add these lines to the “declaration” section:

```
#ifdef HAVE_CHARSET_MYSET
extern CHARSET_INFO my_charset_MYSET_general_ci;
extern CHARSET_INFO my_charset_MYSET_bin;
#endif
```

   Add these lines to the “registration” section:

```
#ifdef HAVE_CHARSET_MYSET
  add_compiled_collation(&my_charset_MYSET_general_ci);
  add_compiled_collation(&my_charset_MYSET_bin);
#endif
```

b. If the character set uses `ctype-MYSET.c`, edit `strings/CMakeLists.txt` and add `ctype-MYSET.c` to the definition of the `STRINGS_SOURCES` variable.

c. Edit `cmake/character_sets.cmake`:

   i. Add `MYSET` to the value of with `CHARSETS_AVAILABLE` in alphabetic order.

   ii. Add `MYSET` to the value of `CHARSETS_COMPLEX` in alphabetic order. This is needed even for simple character sets, or `CMake` will not recognize `-DDEFAULT_CHARSET=MYSET`.

4. Reconfigure, recompile, and test.

10.3.1 Character Definition Arrays

Each simple character set has a configuration file located in the `sql/share/charsets` directory. For a character set named `MYSYS`, the file is named `MYSET.xml`. It uses `<map>` array elements to list character set properties. `<map>` elements appear within these elements:

- `<ctype>` defines attributes for each character.
- `<lower>` and `<upper>` list the lowercase and uppercase characters.
- `<unicode>` maps 8-bit character values to Unicode values.
- `<collation>` elements indicate character ordering for comparison and sorting, one element per collation. Binary collations need no `<map>` element because the character codes themselves provide the ordering.

For a complex character set as implemented in a `ctype-MYSET.c` file in the `strings` directory, there are corresponding arrays: `ctype_MYSET[]`, `to_lower_MYSET[]`, and so forth. Not every complex character set has all of the arrays. See also the existing `ctype-*.c` files for examples. See the `CHARSET_INFO.txt` file in the `strings` directory for additional information.

Most of the arrays are indexed by character value and have 256 elements. The `<ctype>` array is indexed by character value + 1 and has 257 elements. This is a legacy convention for handling `EOF`.

`<ctype>` array elements are bit values. Each element describes the attributes of a single character in the character set. Each attribute is associated with a bitmask, as defined in `include/m_ctype.h`:

```
#define _MY_U   01      /* Upper case */
#define _MY_L   02      /* Lower case */
#define _MY_NMR 04      /* Numeral (digit) */
#define _MY_SPC 010     /* Spacing character */
#define _MY_PNT 020     /* Punctuation */
#define _MY_CTR 040     /* Control character */
#define _MY_B   0100    /* Blank */
```

```
#define _MY_X   0200    /* heXadecimal digit */
```

The <ctype> value for a given character should be the union of the applicable bitmask values that describe the character. For example, 'A' is an uppercase character (_MY_U) as well as a hexadecimal digit (_MY_X), so its ctype value should be defined like this:

```
ctype['A'+1] = _MY_U | _MY_X = 01 | 0200 = 0201
```

The bitmask values in m_ctype.h are octal values, but the elements of the <ctype> array in *MYSET*.xml should be written as hexadecimal values.

The <lower> and <upper> arrays hold the lowercase and uppercase characters corresponding to each member of the character set. For example:

```
lower['A'] should contain 'a'
upper['a'] should contain 'A'
```

Each <collation> array indicates how characters should be ordered for comparison and sorting purposes. MySQL sorts characters based on the values of this information. In some cases, this is the same as the <upper> array, which means that sorting is case-insensitive. For more complicated sorting rules (for complex character sets), see the discussion of string collating in Section 10.3.2, “String Collating Support for Complex Character Sets”.

10.3.2 String Collating Support for Complex Character Sets

For a simple character set named *MYSET*, sorting rules are specified in the *MYSET*.xml configuration file using <map> array elements within <collation> elements. If the sorting rules for your language are too complex to be handled with simple arrays, you must define string collating functions in the ctype-*MYSET*.c source file in the strings directory.

The existing character sets provide the best documentation and examples to show how these functions are implemented. Look at the ctype-*.c files in the strings directory, such as the files for the big5, czech, gbk, sjis, and tis160 character sets. Take a look at the MY_COLLATION_HANDLER structures to see how they are used. See also the CHARSET_INFO.txt file in the strings directory for additional information.

10.3.3 Multi-Byte Character Support for Complex Character Sets

If you want to add support for a new character set named *MYSET* that includes multibyte characters, you must use multibyte character functions in the ctype-*MYSET*.c source file in the strings directory.

The existing character sets provide the best documentation and examples to show how these functions are implemented. Look at the ctype-*.c files in the strings directory, such as the files for the euc_kr, gb2312, gbk, sjis, and ujis character sets. Take a look at the MY_CHARSET_HANDLER structures to see how they are used. See also the CHARSET_INFO.txt file in the strings directory for additional information.

10.4 Adding a Collation to a Character Set

A collation is a set of rules that defines how to compare and sort character strings. Each collation in MySQL belongs to a single character set. Every character set has at least one collation, and most have two or more collations.

A collation orders characters based on weights. Each character in a character set maps to a weight. Characters with equal weights compare as equal, and characters with unequal weights compare according to the relative magnitude of their weights.

The WEIGHT_STRING() function can be used to see the weights for the characters in a string. The value that it returns to indicate weights is a binary string, so it is convenient to use

`HEX(WEIGHT_STRING(str))` to display the weights in printable form. The following example shows that weights do not differ for lettercase for the letters in `'AaBb'` if it is a nonbinary case-insensitive string, but do differ if it is a binary string:

```
mysql> SELECT HEX(WEIGHT_STRING('AaBb' COLLATE latin1_swedish_ci));
+-----------------------------------------------------+
| HEX(WEIGHT_STRING('AaBb' COLLATE latin1_swedish_ci)) |
+-----------------------------------------------------+
| 41414242                                            |
+-----------------------------------------------------+
mysql> SELECT HEX(WEIGHT_STRING(BINARY 'AaBb'));
+----------------------------------+
| HEX(WEIGHT_STRING(BINARY 'AaBb')) |
+----------------------------------+
| 41614262                         |
+----------------------------------+
```

MySQL supports several collation implementations, as discussed in Section 10.4.1, “Collation Implementation Types”. Some of these can be added to MySQL without recompiling:

- Simple collations for 8-bit character sets.
- UCA-based collations for Unicode character sets.
- Binary (*xxx*`_bin`) collations.

The following sections describe how to add collations of the first two types to existing character sets. All existing character sets already have a binary collation, so there is no need here to describe how to add one.

Summary of the procedure for adding a new collation:

1. Choose a collation ID.
2. Add configuration information that names the collation and describes the character-ordering rules.
3. Restart the server.
4. Verify that the collation is present.

The instructions here cover only collations that can be added without recompiling MySQL. To add a collation that does require recompiling (as implemented by means of functions in a C source file), use the instructions in Section 10.3, “Adding a Character Set”. However, instead of adding all the information required for a complete character set, just modify the appropriate files for an existing character set. That is, based on what is already present for the character set's current collations, add data structures, functions, and configuration information for the new collation.

Note

If you modify an existing collation, that may affect the ordering of rows for indexes on columns that use the collation. In this case, rebuild any such indexes to avoid problems such as incorrect query results. See Section 2.10.3, “Rebuilding or Repairing Tables or Indexes”.

Additional Resources

- The Unicode Collation Algorithm (UCA) specification: http://www.unicode.org/reports/tr10/
- The Locale Data Markup Language (LDML) specification: http://www.unicode.org/reports/tr35/

10.4.1 Collation Implementation Types

MySQL implements several types of collations:

Simple collations for 8-bit character sets

This kind of collation is implemented using an array of 256 weights that defines a one-to-one mapping from character codes to weights. `latin1_swedish_ci` is an example. It is a case-insensitive collation, so the uppercase and lowercase versions of a character have the same weights and they compare as equal.

```
mysql> SET NAMES 'latin1' COLLATE 'latin1_swedish_ci';
Query OK, 0 rows affected (0.01 sec)

mysql> SELECT HEX(WEIGHT_STRING('a')), HEX(WEIGHT_STRING('A'));
+-------------------------+-------------------------+
| HEX(WEIGHT_STRING('a')) | HEX(WEIGHT_STRING('A')) |
+-------------------------+-------------------------+
| 41                      | 41                      |
+-------------------------+-------------------------+
1 row in set (0.01 sec)

mysql> SELECT 'a' = 'A';
+-----------+
| 'a' = 'A' |
+-----------+
|         1 |
+-----------+
1 row in set (0.12 sec)
```

For implementation instructions, see Section 10.4.3, “Adding a Simple Collation to an 8-Bit Character Set”.

Complex collations for 8-bit character sets

This kind of collation is implemented using functions in a C source file that define how to order characters, as described in Section 10.3, “Adding a Character Set”.

Collations for non-Unicode multibyte character sets

For this type of collation, 8-bit (single-byte) and multibyte characters are handled differently. For 8-bit characters, character codes map to weights in case-insensitive fashion. (For example, the single-byte characters `'a'` and `'A'` both have a weight of `0x41`.) For multibyte characters, there are two types of relationship between character codes and weights:

- Weights equal character codes. `sjis_japanese_ci` is an example of this kind of collation. The multibyte character `'ぢ'` has a character code of `0x82C0`, and the weight is also `0x82C0`.

```
mysql> CREATE TABLE t1
       (c1 VARCHAR(2) CHARACTER SET sjis COLLATE sjis_japanese_ci);
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO t1 VALUES ('a'),('A'),(0x82C0);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT c1, HEX(c1), HEX(WEIGHT_STRING(c1)) FROM t1;
+------+---------+------------------------+
| c1   | HEX(c1) | HEX(WEIGHT_STRING(c1)) |
+------+---------+------------------------+
| a    | 61      | 41                     |
| A    | 41      | 41                     |
| ぢ    | 82C0    | 82C0                   |
+------+---------+------------------------+
3 rows in set (0.00 sec)
```

- Character codes map one-to-one to weights, but a code is not necessarily equal to the weight. `gbk_chinese_ci` is an example of this kind of collation. The multibyte character `'膰'` has a character code of `0x81B0` but a weight of `0xC286`.

```
mysql> CREATE TABLE t1
       (c1 VARCHAR(2) CHARACTER SET gbk COLLATE gbk_chinese_ci);
Query OK, 0 rows affected (0.33 sec)

mysql> INSERT INTO t1 VALUES ('a'),('A'),(0x81B0);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT c1, HEX(c1), HEX(WEIGHT_STRING(c1)) FROM t1;
+------+---------+-----------------------+
| c1   | HEX(c1) | HEX(WEIGHT_STRING(c1)) |
+------+---------+-----------------------+
| a    | 61      | 41                     |
| A    | 41      | 41                     |
| 膰    | 81B0    | C286                   |
+------+---------+-----------------------+
3 rows in set (0.00 sec)
```

For implementation instructions, see Section 10.3, “Adding a Character Set”.

Collations for Unicode multibyte character sets

Some of these collations are based on the Unicode Collation Algorithm (UCA), others are not.

Non-UCA collations have a one-to-one mapping from character code to weight. In MySQL, such collations are case insensitive and accent insensitive. `utf8_general_ci` is an example: `'a'`, `'A'`, `'À'`, and `'á'` each have different character codes but all have a weight of `0x0041` and compare as equal.

```
mysql> SET NAMES 'utf8' COLLATE 'utf8_general_ci';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t1
       (c1 CHAR(1) CHARACTER SET UTF8 COLLATE utf8_general_ci);
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO t1 VALUES ('a'),('A'),('À'),('á');
Query OK, 4 rows affected (0.00 sec)
Records: 4  Duplicates: 0  Warnings: 0

mysql> SELECT c1, HEX(c1), HEX(WEIGHT_STRING(c1)) FROM t1;
+------+---------+-----------------------+
| c1   | HEX(c1) | HEX(WEIGHT_STRING(c1)) |
+------+---------+-----------------------+
| a    | 61      | 0041                   |
| A    | 41      | 0041                   |
| À    | C380    | 0041                   |
| á    | C3A1    | 0041                   |
+------+---------+-----------------------+
4 rows in set (0.00 sec)
```

UCA-based collations in MySQL have these properties:

- If a character has weights, each weight uses 2 bytes (16 bits).
- A character may have zero weights (or an empty weight). In this case, the character is ignorable. Example: "U+0000 NULL" does not have a weight and is ignorable.
- A character may have one weight. Example: `'a'` has a weight of `0x0E33`.

```
mysql> SET NAMES 'utf8' COLLATE 'utf8_unicode_ci';
Query OK, 0 rows affected (0.05 sec)

mysql> SELECT HEX('a'), HEX(WEIGHT_STRING('a'));
+----------+------------------------+
```

```
| HEX('a') | HEX(WEIGHT_STRING('a')) |
+----------+-------------------------+
| 61       | 0E33                    |
+----------+-------------------------+
1 row in set (0.02 sec)
```

- A character may have many weights. This is an expansion. Example: The German letter `'ß'` (SZ ligature, or SHARP S) has a weight of `0x0FEA0FEA`.

```
mysql> SET NAMES 'utf8' COLLATE 'utf8_unicode_ci';
Query OK, 0 rows affected (0.11 sec)

mysql> SELECT HEX('ß'), HEX(WEIGHT_STRING('ß'));
+-----------+--------------------------+
| HEX('ß')  | HEX(WEIGHT_STRING('ß'))  |
+-----------+--------------------------+
| C39F      | 0FEA0FEA                 |
+-----------+--------------------------+
1 row in set (0.00 sec)
```

- Many characters may have one weight. This is a contraction. Example: `'ch'` is a single letter in Czech and has a weight of `0x0EE2`.

```
mysql> SET NAMES 'utf8' COLLATE 'utf8_czech_ci';
Query OK, 0 rows affected (0.09 sec)

mysql> SELECT HEX('ch'), HEX(WEIGHT_STRING('ch'));
+-----------+--------------------------+
| HEX('ch') | HEX(WEIGHT_STRING('ch')) |
+-----------+--------------------------+
| 6368      | 0EE2                     |
+-----------+--------------------------+
1 row in set (0.00 sec)
```

A many-characters-to-many-weights mapping is also possible (this is contraction with expansion), but is not supported by MySQL.

For implementation instructions, for a non-UCA collation, see Section 10.3, “Adding a Character Set”. For a UCA collation, see Section 10.4.4, “Adding a UCA Collation to a Unicode Character Set”.

Miscellaneous collations

There are also a few collations that do not fall into any of the previous categories.

10.4.2 Choosing a Collation ID

Each collation must have a unique ID. To add a collation, you must choose an ID value that is not currently used. MySQL supports two-byte collation IDs. The range of IDs from 1024 to 2047 is reserved for user-defined collations. The collation ID that you choose will appear in these contexts:

- The `ID` column of the `INFORMATION_SCHEMA.COLLATIONS` table.
- The `Id` column of `SHOW COLLATION` output.
- The `charsetnr` member of the `MYSQL_FIELD` C API data structure.
- The `number` member of the `MY_CHARSET_INFO` data structure returned by the `mysql_get_character_set_info()` C API function.

To determine the largest currently used ID, issue the following statement:

```
mysql> SELECT MAX(ID) FROM INFORMATION_SCHEMA.COLLATIONS;
+---------+
```

```
| MAX(ID) |
+---------+
|     210 |
+---------+
```

To display a list of all currently used IDs, issue this statement:

```
mysql> SELECT ID FROM INFORMATION_SCHEMA.COLLATIONS ORDER BY ID;
+-----+
| ID  |
+-----+
|   1 |
|   2 |
| ... |
|  52 |
|  53 |
|  57 |
|  58 |
| ... |
|  98 |
|  99 |
| 128 |
| 129 |
| ... |
| 210 |
+-----+
```

Warning

Before MySQL 5.5, which provides for a range of user-defined collation IDs, you must choose an ID in the range from 1 to 254. In this case, if you upgrade MySQL, you may find that the collation ID you choose has been assigned to a collation included in the new MySQL distribution. In this case, you will need to choose a new value for your own collation.

In addition, before upgrading, you should save the configuration files that you change. If you upgrade in place, the process will replace the your modified files.

10.4.3 Adding a Simple Collation to an 8-Bit Character Set

This section describes how to add a simple collation for an 8-bit character set by writing the `<collation>` elements associated with a `<charset>` character set description in the MySQL `Index.xml` file. The procedure described here does not require recompiling MySQL. The example adds a collation named `latin1_test_ci` to the `latin1` character set.

1. Choose a collation ID, as shown in Section 10.4.2, “Choosing a Collation ID”. The following steps use an ID of 1024.

2. Modify the `Index.xml` and `latin1.xml` configuration files. These files are located in the directory named by the `character_sets_dir` system variable. You can check the variable value as follows, although the path name might be different on your system:

   ```
   mysql> SHOW VARIABLES LIKE 'character_sets_dir';
   +--------------------+----------------------------------------+
   | Variable_name      | Value                                  |
   +--------------------+----------------------------------------+
   | character_sets_dir | /user/local/mysql/share/mysql/charsets/ |
   +--------------------+----------------------------------------+
   ```

3. Choose a name for the collation and list it in the `Index.xml` file. Find the `<charset>` element for the character set to which the collation is being added, and add a `<collation>` element that indicates the collation name and ID, to associate the name with the ID. For example:

```
<charset name="latin1">
  ...
  <collation name="latin1_test_ci" id="1024"/>
  ...
</charset>
```

4. In the `latin1.xml` configuration file, add a `<collation>` element that names the collation and that contains a `<map>` element that defines a character code-to-weight mapping table for character codes 0 to 255. Each value within the `<map>` element must be a number in hexadecimal format.

```
<collation name="latin1_test_ci">
<map>
 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
 10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F
 20 21 22 23 24 25 26 27 28 29 2A 2B 2C 2D 2E 2F
 30 31 32 33 34 35 36 37 38 39 3A 3B 3C 3D 3E 3F
 40 41 42 43 44 45 46 47 48 49 4A 4B 4C 4D 4E 4F
 50 51 52 53 54 55 56 57 58 59 5A 5B 5C 5D 5E 5F
 60 41 42 43 44 45 46 47 48 49 4A 4B 4C 4D 4E 4F
 50 51 52 53 54 55 56 57 58 59 5A 7B 7C 7D 7E 7F
 80 81 82 83 84 85 86 87 88 89 8A 8B 8C 8D 8E 8F
 90 91 92 93 94 95 96 97 98 99 9A 9B 9C 9D 9E 9F
 A0 A1 A2 A3 A4 A5 A6 A7 A8 A9 AA AB AC AD AE AF
 B0 B1 B2 B3 B4 B5 B6 B7 B8 B9 BA BB BC BD BE BF
 41 41 41 41 5B 5D 5B 43 45 45 45 45 49 49 49 49
 44 4E 4F 4F 4F 4F 5C D7 5C 55 55 55 59 59 DE DF
 41 41 41 41 5B 5D 5B 43 45 45 45 45 49 49 49 49
 44 4E 4F 4F 4F 4F 5C F7 5C 55 55 55 59 59 DE FF
</map>
</collation>
```

5. Restart the server and use this statement to verify that the collation is present:

```
mysql> SHOW COLLATION WHERE Collation = 'latin1_test_ci';
+----------------+---------+------+---------+----------+---------+
| Collation      | Charset | Id   | Default | Compiled | Sortlen |
+----------------+---------+------+---------+----------+---------+
| latin1_test_ci | latin1  | 1024 |         |          |       1 |
+----------------+---------+------+---------+----------+---------+
```

10.4.4 Adding a UCA Collation to a Unicode Character Set

This section describes how to add a UCA collation for a Unicode character set by writing the `<collation>` element within a `<charset>` character set description in the MySQL `Index.xml` file. The procedure described here does not require recompiling MySQL. It uses a subset of the Locale Data Markup Language (LDML) specification, which is available at http://www.unicode.org/reports/tr35/. With this method, you need not define the entire collation. Instead, you begin with an existing “base” collation and describe the new collation in terms of how it differs from the base collation. The following table lists the base collations of the Unicode character sets for which UCA collations can be defined. It is not possible to create user-defined UCA collations for `utf16le`; there is no `utf16le_unicode_ci` collation that would serve as the basis for such collations.

Table 10.4 MySQL Character Sets Available for User-Defined UCA Collations

Character Set	Base Collation
`utf8`	`utf8_unicode_ci`
`ucs2`	`ucs2_unicode_ci`
`utf16`	`utf16_unicode_ci`
`utf32`	`utf32_unicode_ci`

The following sections show how to add a collation that is defined using LDML syntax, and provide a summary of LDML rules supported in MySQL.

10.4.4.1 Defining a UCA Collation Using LDML Syntax

To add a UCA collation for a Unicode character set without recompiling MySQL, use the following procedure. If you are unfamiliar with the LDML rules used to describe the collation's sort characteristics, see Section 10.4.4.2, “LDML Syntax Supported in MySQL”.

The example adds a collation named `utf8_phone_ci` to the `utf8` character set. The collation is designed for a scenario involving a Web application for which users post their names and phone numbers. Phone numbers can be given in very different formats:

```
+7-12345-67
+7-12-345-67
+7 12 345 67
+7 (12) 345 67
+71234567
```

The problem raised by dealing with these kinds of values is that the varying permissible formats make searching for a specific phone number very difficult. The solution is to define a new collation that reorders punctuation characters, making them ignorable.

1. Choose a collation ID, as shown in Section 10.4.2, “Choosing a Collation ID”. The following steps use an ID of 1029.

2. To modify the `Index.xml` configuration file. This file is located in the directory named by the `character_sets_dir` system variable. You can check the variable value as follows, although the path name might be different on your system:

   ```
   mysql> SHOW VARIABLES LIKE 'character_sets_dir';
   +--------------------+----------------------------------------+
   | Variable_name      | Value                                  |
   +--------------------+----------------------------------------+
   | character_sets_dir | /user/local/mysql/share/mysql/charsets/ |
   +--------------------+----------------------------------------+
   ```

3. Choose a name for the collation and list it in the `Index.xml` file. In addition, you'll need to provide the collation ordering rules. Find the `<charset>` element for the character set to which the collation is being added, and add a `<collation>` element that indicates the collation name and ID, to associate the name with the ID. Within the `<collation>` element, provide a `<rules>` element containing the ordering rules:

   ```
   <charset name="utf8">
     ...
     <collation name="utf8_phone_ci" id="1029">
       <rules>
         <reset>\u0000</reset>
         <i>\u0020</i> <!-- space -->
         <i>\u0028</i> <!-- left parenthesis -->
         <i>\u0029</i> <!-- right parenthesis -->
         <i>\u002B</i> <!-- plus -->
         <i>\u002D</i> <!-- hyphen -->
       </rules>
     </collation>
     ...
   </charset>
   ```

4. If you want a similar collation for other Unicode character sets, add other `<collation>` elements. For example, to define `ucs2_phone_ci`, add a `<collation>` element to the `<charset name="ucs2">` element. Remember that each collation must have its own unique ID.

5. Restart the server and use this statement to verify that the collation is present:

   ```
   mysql> SHOW COLLATION WHERE Collation = 'utf8_phone_ci';
   ```

```
+---------------+---------+------+---------+----------+---------+
| Collation     | Charset | Id   | Default | Compiled | Sortlen |
+---------------+---------+------+---------+----------+---------+
| utf8_phone_ci | utf8    | 1029 |         |          |       8 |
+---------------+---------+------+---------+----------+---------+
```

Now test the collation to make sure that it has the desired properties.

Create a table containing some sample phone numbers using the new collation:

```
mysql> CREATE TABLE phonebook (
         name VARCHAR(64),
         phone VARCHAR(64) CHARACTER SET utf8 COLLATE utf8_phone_ci
       );
Query OK, 0 rows affected (0.09 sec)

mysql> INSERT INTO phonebook VALUES ('Svoj','+7 912 800 80 02');
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO phonebook VALUES ('Hf','+7 (912) 800 80 04');
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO phonebook VALUES ('Bar','+7-912-800-80-01');
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO phonebook VALUES ('Ramil','(7912) 800 80 03');
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO phonebook VALUES ('Sanja','+380 (912) 8008005');
Query OK, 1 row affected (0.00 sec)
```

Run some queries to see whether the ignored punctuation characters are in fact ignored for comparison and sorting:

```
mysql> SELECT * FROM phonebook ORDER BY phone;
+-------+--------------------+
| name  | phone              |
+-------+--------------------+
| Sanja | +380 (912) 8008005 |
| Bar   | +7-912-800-80-01   |
| Svoj  | +7 912 800 80 02   |
| Ramil | (7912) 800 80 03   |
| Hf    | +7 (912) 800 80 04 |
+-------+--------------------+
5 rows in set (0.00 sec)

mysql> SELECT * FROM phonebook WHERE phone='+7(912)800-80-01';
+------+------------------+
| name | phone            |
+------+------------------+
| Bar  | +7-912-800-80-01 |
+------+------------------+
1 row in set (0.00 sec)

mysql> SELECT * FROM phonebook WHERE phone='79128008001';
+------+------------------+
| name | phone            |
+------+------------------+
| Bar  | +7-912-800-80-01 |
+------+------------------+
1 row in set (0.00 sec)

mysql> SELECT * FROM phonebook WHERE phone='7 9 1 2 8 0 0 8 0 0 1';
+------+------------------+
| name | phone            |
+------+------------------+
| Bar  | +7-912-800-80-01 |
+------+------------------+
1 row in set (0.00 sec)
```

10.4.4.2 LDML Syntax Supported in MySQL

This section describes the LDML syntax that MySQL recognizes. This is a subset of the syntax described in the LDML specification available at http://www.unicode.org/reports/tr35/, which should be consulted for further information. MySQL recognizes a large enough subset of the syntax that, in many cases, it is possible to download a collation definition from the Unicode Common Locale Data Repository and paste the relevant part (that is, the part between the `<rules>` and `</rules>` tags) into the MySQL `Index.xml` file. The rules described here are all supported except that character sorting occurs only at the primary level. Rules that specify differences at secondary or higher sort levels are recognized (and thus can be included in collation definitions) but are treated as equality at the primary level.

The MySQL server generates diagnostics when it finds problems while parsing the `Index.xml` file. See Section 10.4.4.3, “Diagnostics During Index.xml Parsing”.

Character Representation

Characters named in LDML rules can be written literally or in `\u`*nnnn* format, where *nnnn* is the hexadecimal Unicode code point value. For example, `A` and `á` can be written literally or as `\u0041` and `\u00E1`. Within hexadecimal values, the digits `A` through `F` are not case sensitive; `\u00E1` and `\u00e1` are equivalent. For UCA 4.0.0 collations, hexadecimal notation can be used only for characters in the Basic Multilingual Plane, not for characters outside the BMP range of `0000` to `FFFF`. For UCA 5.2.0 collations, hexadecimal notation can be used for any character.

The `Index.xml` file itself should be written using UTF-8 encoding.

Syntax Rules

LDML has reset rules and shift rules to specify character ordering. Orderings are given as a set of rules that begin with a reset rule that establishes an anchor point, followed by shift rules that indicate how characters sort relative to the anchor point.

- A `<reset>` rule does not specify any ordering in and of itself. Instead, it “resets” the ordering for subsequent shift rules to cause them to be taken in relation to a given character. Either of the following rules resets subsequent shift rules to be taken in relation to the letter `'A'`:

  ```
  <reset>A</reset>

  <reset>\u0041</reset>
  ```

- The `<p>`, `<s>`, and `<t>` shift rules define primary, secondary, and tertiary differences of a character from another character:

  - Use primary differences to distinguish separate letters.
  - Use secondary differences to distinguish accent variations.
  - Use tertiary differences to distinguish lettercase variations.

  Either of these rules specifies a primary shift rule for the `'G'` character:

  ```
  <p>G</p>

  <p>\u0047</p>
  ```

- The `<i>` shift rule indicates that one character sorts identically to another. The following rules cause `'b'` to sort the same as `'a'`:

  ```
  <reset>a</reset>
  <i>b</i>
  ```

- Abbreviated shift syntax specifies multiple shift rules using a single pair of tags. The following table shows the correspondence between abbreviated syntax rules and the equivalent nonabbreviated rules.

  Table 10.5 Abbreviated Shift Syntax

Abbreviated Syntax	Nonabbreviated Syntax
`<pc>xyz</pc>`	`<p>x</p><p>y</p><p>z</p>`
`<sc>xyz</sc>`	`<s>x</s><s>y</s><s>z</s>`
`<tc>xyz</tc>`	`<t>x</t><t>y</t><t>z</t>`
`<ic>xyz</ic>`	`<i>x</i><i>y</i><i>z</i>`

- An expansion is a reset rule that establishes an anchor point for a multiple-character sequence. MySQL supports expansions 2 to 6 characters long. The following rules put `'z'` greater at the primary level than the sequence of three characters `'abc'`:

```
<reset>abc</reset>
<p>z</p>
```

- A contraction is a shift rule that sorts a multiple-character sequence. MySQL supports contractions 2 to 6 characters long. The following rules put the sequence of three characters `'xyz'` greater at the primary level than `'a'`:

```
<reset>a</reset>
<p>xyz</p>
```

- Long expansions and long contractions can be used together. These rules put the sequence of three characters `'xyz'` greater at the primary level than the sequence of three characters `'abc'`:

```
<reset>abc</reset>
<p>xyz</p>
```

- Normal expansion syntax uses `<x>` plus `<extend>` elements to specify an expansion. The following rules put the character `'k'` greater at the secondary level than the sequence `'ch'`. That is, `'k'` behaves as if it expands to a character after `'c'` followed by `'h'`:

```
<reset>c</reset>
<x><s>k</s><extend>h</extend></x>
```

  This syntax permits long sequences. These rules sort the sequence `'ccs'` greater at the tertiary level than the sequence `'cscs'`:

```
<reset>cs</reset>
<x><t>ccs</t><extend>cs</extend></x>
```

  The LDML specification describes normal expansion syntax as “tricky.” See that specification for details.

- Previous context syntax uses `<x>` plus `<context>` elements to specify that the context before a character affects how it sorts. The following rules put `'-'` greater at the secondary level than `'a'`, but only when `'-'` occurs after `'b'`:

```
<reset>a</reset>
<x><context>b</context><s>-</s></x>
```

- Previous context syntax can include the `<extend>` element. These rules put `'def'` greater at the primary level than `'aghi'`, but only when `'def'` comes after `'abc'`:

```
<reset>a</reset>
<x><context>abc</context><p>def</p><extend>ghi</extend></x>
```

- Reset rules permit a `before` attribute. Normally, shift rules after a reset rule indicate characters that sort after the reset character. Shift rules after a reset rule that has the `before` attribute indicate characters that sort before the reset character. The following rules put the character `'b'` immediately before `'a'` at the primary level:

```
<reset before="primary">a</reset>
<p>b</p>
```

Permissible `before` attribute values specify the sort level by name or the equivalent numeric value:

```
<reset before="primary">
<reset before="1">

<reset before="secondary">
<reset before="2">

<reset before="tertiary">
<reset before="3">
```

- A reset rule can name a logical reset position rather than a literal character:

```
<first_tertiary_ignorable/>
<last_tertiary_ignorable/>
<first_secondary_ignorable/>
<last_secondary_ignorable/>
<first_primary_ignorable/>
<last_primary_ignorable/>
<first_variable/>
<last_variable/>
<first_non_ignorable/>
<last_non_ignorable/>
<first_trailing/>
<last_trailing/>
```

These rules put `'z'` greater at the primary level than nonignorable characters that have a Default Unicode Collation Element Table (DUCET) entry and that are not CJK:

```
<reset><last_non_ignorable/></reset>
<p>z</p>
```

Logical positions have the code points shown in the following table.

Table 10.6 Logical Reset Position Code Points

Logical Position	Unicode 4.0.0 Code Point	Unicode 5.2.0 Code Point
`<first_non_ignorable/>`	U+02D0	U+02D0
`<last_non_ignorable/>`	U+A48C	U+1342E
`<first_primary_ignorable/>`	U+0332	U+0332
`<last_primary_ignorable/>`	U+20EA	U+101FD
`<first_secondary_ignorable/>`	U+0000	U+0000
`<last_secondary_ignorable/>`	U+FE73	U+FE73
`<first_tertiary_ignorable/>`	U+0000	U+0000
`<last_tertiary_ignorable/>`	U+FE73	U+FE73
`<first_trailing/>`	U+0000	U+0000

Logical Position	Unicode 4.0.0 Code Point	Unicode 5.2.0 Code Point
`<last_trailing/>`	U+0000	U+0000
`<first_variable/>`	U+0009	U+0009
`<last_variable/>`	U+2183	U+1D371

- The `<collation>` element permits a `shift-after-method` attribute that affects character weight calculation for shift rules. The attribute has these permitted values:

  - `simple`: Calculate character weights as for reset rules that do not have a `before` attribute. This is the default if the attribute is not given.
  - `expand`: Use expansions for shifts after reset rules.

  Suppose that `'0'` and `'1'` have weights of `0E29` and `0E2A` and we want to put all basic Latin letters between `'0'` and `'1'`:

```
<reset>0</reset>
<pc>abcdefghijklmnopqrstuvwxyz</pc>
```

  For simple shift mode, weights are calculated as follows:

```
'a' has weight 0E29+1
'b' has weight 0E29+2
'c' has weight 0E29+3
...
```

  However, there are not enough vacant positions to put 26 characters between `'0'` and `'1'`. The result is that digits and letters are intermixed.

  To solve this, use `shift-after-method="expand"`. Then weights are calculated like this:

```
'a' has weight [0E29][233D+1]
'b' has weight [0E29][233D+2]
'c' has weight [0E29][233D+3]
...
```

  `233D` is the UCA 4.0.0 weight for character `0xA48C`, which is the last nonignorable character (a sort of the greatest character in the collation, excluding CJK). UCA 5.2.0 is similar but uses `3ACA`, for character `0x1342E`.

MySQL-Specific LDML Extensions

An extension to LDML rules permits the `<collation>` element to include an optional `version` attribute in `<collation>` tags to indicate the UCA version on which the collation is based. If the `version` attribute is omitted, its default value is `4.0.0`. For example, this specification indicates a collation that is based on UCA 5.2.0:

```
<collation id="nnn" name="utf8_xxx_ci" version="5.2.0">
...
</collation>
```

10.4.4.3 Diagnostics During Index.xml Parsing

The MySQL server generates diagnostics when it finds problems while parsing the `Index.xml` file:

- Unknown tags are written to the error log. For example, the following message results if a collation definition contains a `<aaa>` tag:

```
[Warning] Buffered warning: Unknown LDML tag:
'charsets/charset/collation/rules/aaa'
```

- If collation initialization is not possible, the server reports an “Unknown collation” error, and also generates warnings explaining the problems, such as in the previous example. In other cases, when a collation description is generally correct but contains some unknown tags, the collation is initialized and is available for use. The unknown parts are ignored, but a warning is generated in the error log.

- Problems with collations generate warnings that clients can display with `SHOW WARNINGS`. Suppose that a reset rule contains an expansion longer than the maximum supported length of 6 characters:

  ```
  <reset>abcdefghi</reset>
  <i>x</i>
  ```

  An attempt to use the collation produces warnings:

  ```
  mysql> SELECT _utf8'test' COLLATE utf8_test_ci;
  ERROR 1273 (HY000): Unknown collation: 'utf8_test_ci'
  mysql> SHOW WARNINGS;
  +---------+------+----------------------------------------+
  | Level   | Code | Message                                |
  +---------+------+----------------------------------------+
  | Error   | 1273 | Unknown collation: 'utf8_test_ci'      |
  | Warning | 1273 | Expansion is too long at 'abcdefghi=x' |
  +---------+------+----------------------------------------+
  ```

10.5 Character Set Configuration

You can change the default server character set and collation with the `--character-set-server` and `--collation-server` options when you start the server. The collation must be a legal collation for the default character set. (Use the `SHOW COLLATION` statement to determine which collations are available for each character set.) See Section 5.1.4, “Server Command Options”.

If you try to use a character set that is not compiled into your binary, you might run into the following problems:

- Your program uses an incorrect path to determine where the character sets are stored (which is typically the `share/mysql/charsets` or `share/charsets` directory under the MySQL installation directory). This can be fixed by using the `--character-sets-dir` option when you run the program in question. For example, to specify a directory to be used by MySQL client programs, list it in the `[client]` group of your option file. The examples given here show what the setting might look like for Unix or Windows, respectively:

  ```
  [client]
  character-sets-dir=/usr/local/mysql/share/mysql/charsets

  [client]
  character-sets-dir="C:/Program Files/MySQL/MySQL Server 8.0/share/charsets"
  ```

- The character set is a complex character set that cannot be loaded dynamically. In this case, you must recompile the program with support for the character set.

  For Unicode character sets, you can define collations without recompiling by using LDML notation. See Section 10.4.4, “Adding a UCA Collation to a Unicode Character Set”.

- The character set is a dynamic character set, but you do not have a configuration file for it. In this case, you should install the configuration file for the character set from a new MySQL distribution.

- If your character set index file does not contain the name for the character set, your program displays an error message. The file is named `Index.xml` and the message is:

  ```
  ```

```
Character set 'charset_name' is not a compiled character set and is not
specified in the '/usr/share/mysql/charsets/Index.xml' file
```

To solve this problem, you should either get a new index file or manually add the name of any missing character sets to the current file.

You can force client programs to use specific character set as follows:

```
[client]
default-character-set=charset_name
```

This is normally unnecessary. However, when `character_set_system` differs from `character_set_server` or `character_set_client`, and you input characters manually (as database object identifiers, column values, or both), these may be displayed incorrectly in output from the client or the output itself may be formatted incorrectly. In such cases, starting the mysql client with `--default-character-set=system_character_set`—that is, setting the client character set to match the system character set—should fix the problem.

For `MyISAM` tables, you can check the character set name and number for a table with `myisamchk -dvv tbl_name`.

10.6 MySQL Server Time Zone Support

MySQL Server maintains several time zone settings:

- The system time zone. When the server starts, it attempts to determine the time zone of the host machine and uses it to set the `system_time_zone` system variable. The value does not change thereafter.

  You can set the system time zone for MySQL Server at startup with the `--timezone=timezone_name` option to `mysqld_safe`. You can also set it by setting the `TZ` environment variable before you start `mysqld`. The permissible values for `--timezone` or `TZ` are system dependent. Consult your operating system documentation to see what values are acceptable.

- The server's current time zone. The global `time_zone` system variable indicates the time zone the server currently is operating in. The initial value for `time_zone` is `'SYSTEM'`, which indicates that the server time zone is the same as the system time zone.

  The initial global server time zone value can be specified explicitly at startup with the `--default-time-zone=timezone` option on the command line, or you can use the following line in an option file:

  ```
  default-time-zone='timezone'
  ```

  If you have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege, you can set the global server time zone value at runtime with this statement:

  ```
  mysql> SET GLOBAL time_zone = timezone;
  ```

- Per-connection time zones. Each client that connects has its own time zone setting, given by the session `time_zone` variable. Initially, the session variable takes its value from the global `time_zone` variable, but the client can change its own time zone with this statement:

  ```
  mysql> SET time_zone = timezone;
  ```

The current session time zone setting affects display and storage of time values that are zone-sensitive. This includes the values displayed by functions such as `NOW()` or `CURTIME()`, and values

stored in and retrieved from `TIMESTAMP` columns. Values for `TIMESTAMP` columns are converted from the current time zone to UTC for storage, and from UTC to the current time zone for retrieval.

The current time zone setting does not affect values displayed by functions such as `UTC_TIMESTAMP()` or values in `DATE`, `TIME`, or `DATETIME` columns. Nor are values in those data types stored in UTC; the time zone applies for them only when converting from `TIMESTAMP` values. If you want locale-specific arithmetic for `DATE`, `TIME`, or `DATETIME` values, convert them to UTC, perform the arithmetic, and then convert back.

The current values of the global and client-specific time zones can be retrieved like this:

```
mysql> SELECT @@global.time_zone, @@session.time_zone;
```

`timezone` values can be given in several formats, none of which are case sensitive:

- The value `'SYSTEM'` indicates that the time zone should be the same as the system time zone.
- The value can be given as a string indicating an offset from UTC, such as `'+10:00'` or `'-6:00'`.
- The value can be given as a named time zone, such as `'Europe/Helsinki'`, `'US/Eastern'`, or `'MET'`. Named time zones can be used only if the time zone information tables in the `mysql` database have been created and populated.

Populating the Time Zone Tables

Several tables in the `mysql` system database exist to maintain time zone information (see Section 5.3, "The mysql System Database"). The MySQL installation procedure creates the time zone tables, but does not load them. You must do so manually using the following instructions.

Note

Loading the time zone information is not necessarily a one-time operation because the information changes occasionally. When such changes occur, applications that use the old rules become out of date and you may find it necessary to reload the time zone tables to keep the information used by your MySQL server current. See the notes at the end of this section.

If your system has its own *zoneinfo* database (the set of files describing time zones), you should use the `mysql_tzinfo_to_sql` program for filling the time zone tables. Examples of such systems are Linux, FreeBSD, Solaris, and OS X. One likely location for these files is the `/usr/share/zoneinfo` directory. If your system does not have a zoneinfo database, you can use the downloadable package described later in this section.

The `mysql_tzinfo_to_sql` program is used to load the time zone tables. On the command line, pass the zoneinfo directory path name to `mysql_tzinfo_to_sql` and send the output into the `mysql` program. For example:

```
shell> mysql_tzinfo_to_sql /usr/share/zoneinfo | mysql -u root mysql
```

`mysql_tzinfo_to_sql` reads your system's time zone files and generates SQL statements from them. `mysql` processes those statements to load the time zone tables.

`mysql_tzinfo_to_sql` also can be used to load a single time zone file or to generate leap second information:

- To load a single time zone file *`tz_file`* that corresponds to a time zone name *`tz_name`*, invoke `mysql_tzinfo_to_sql` like this:

```
shell> mysql_tzinfo_to_sql tz_file tz_name | mysql -u root mysql
```

With this approach, you must execute a separate command to load the time zone file for each named zone that the server needs to know about.

- If your time zone needs to account for leap seconds, initialize the leap second information like this, where *`tz_file`* is the name of your time zone file:

```
shell> mysql_tzinfo_to_sql --leap tz_file | mysql -u root mysql
```

- After running `mysql_tzinfo_to_sql`, it is best to restart the server so that it does not continue to use any previously cached time zone data.

If your system is one that has no zoneinfo database (for example, Windows), you can use a package that is available for download at the MySQL Developer Zone:

```
http://dev.mysql.com/downloads/timezones.html
```

Download a time zone package that contains SQL statements and unpack it, then load the package file contents into the time zone tables:

```
shell> mysql -u root mysql < file_name
```

Then restart the server.

Warning

Do *not* use a downloadable package that contains `MyISAM` tables. MySQL uses `InnoDB` for the time zone tables. Trying to replace them with `MyISAM` tables will cause problems.

Warning

Do not use a downloadable package if your system has a zoneinfo database. Use the `mysql_tzinfo_to_sql` utility instead. Otherwise, you may cause a difference in datetime handling between MySQL and other applications on your system.

For information about time zone settings in replication setup, please see Section 18.4.1, “Replication Features and Issues”.

10.6.1 Staying Current with Time Zone Changes

When time zone rules change, applications that use the old rules become out of date. To stay current, it is necessary to make sure that your system uses current time zone information is used. For MySQL, there are two factors to consider in staying current:

- The operating system time affects the value that the MySQL server uses for times if its time zone is set to `SYSTEM`. Make sure that your operating system is using the latest time zone information. For most operating systems, the latest update or service pack prepares your system for the time changes. Check the Web site for your operating system vendor for an update that addresses the time changes.

- If you replace the system's `/etc/localtime` timezone file with a version that uses rules differing from those in effect at `mysqld` startup, you should restart `mysqld` so that it uses the updated rules. Otherwise, `mysqld` might not notice when the system changes its time.

- If you use named time zones with MySQL, make sure that the time zone tables in the `mysql` database are up to date. If your system has its own zoneinfo database, you should reload the MySQL time zone tables whenever the zoneinfo database is updated. For systems that do not

have their own zoneinfo database, check the MySQL Developer Zone for updates. When a new update is available, download it and use it to replace the content of your current time zone tables. For instructions for both methods, see Populating the Time Zone Tables. `mysqld` caches time zone information that it looks up, so after updating the time zone tables, you should restart `mysqld` to make sure that it does not continue to serve outdated time zone data.

If you are uncertain whether named time zones are available, for use either as the server's time zone setting or by clients that set their own time zone, check whether your time zone tables are empty. The following query determines whether the table that contains time zone names has any rows:

```
mysql> SELECT COUNT(*) FROM mysql.time_zone_name;
+----------+
| COUNT(*) |
+----------+
|        0 |
+----------+
```

A count of zero indicates that the table is empty. In this case, no one can be using named time zones, and you don't need to update the tables. A count greater than zero indicates that the table is not empty and that its contents are available to be used for named time zone support. In this case, you should be sure to reload your time zone tables so that anyone who uses named time zones will get correct query results.

To check whether your MySQL installation is updated properly for a change in Daylight Saving Time rules, use a test like the one following. The example uses values that are appropriate for the 2007 DST 1-hour change that occurs in the United States on March 11 at 2 a.m.

The test uses these two queries:

```
SELECT CONVERT_TZ('2007-03-11 2:00:00','US/Eastern','US/Central');
SELECT CONVERT_TZ('2007-03-11 3:00:00','US/Eastern','US/Central');
```

The two time values indicate the times at which the DST change occurs, and the use of named time zones requires that the time zone tables be used. The desired result is that both queries return the same result (the input time, converted to the equivalent value in the 'US/Central' time zone).

Before updating the time zone tables, you would see an incorrect result like this:

```
mysql> SELECT CONVERT_TZ('2007-03-11 2:00:00','US/Eastern','US/Central');
+-----------------------------------------------------------+
| CONVERT_TZ('2007-03-11 2:00:00','US/Eastern','US/Central') |
+-----------------------------------------------------------+
| 2007-03-11 01:00:00                                       |
+-----------------------------------------------------------+

mysql> SELECT CONVERT_TZ('2007-03-11 3:00:00','US/Eastern','US/Central');
+-----------------------------------------------------------+
| CONVERT_TZ('2007-03-11 3:00:00','US/Eastern','US/Central') |
+-----------------------------------------------------------+
| 2007-03-11 02:00:00                                       |
+-----------------------------------------------------------+
```

After updating the tables, you should see the correct result:

```
mysql> SELECT CONVERT_TZ('2007-03-11 2:00:00','US/Eastern','US/Central');
+-----------------------------------------------------------+
| CONVERT_TZ('2007-03-11 2:00:00','US/Eastern','US/Central') |
+-----------------------------------------------------------+
| 2007-03-11 01:00:00                                       |
+-----------------------------------------------------------+

mysql> SELECT CONVERT_TZ('2007-03-11 3:00:00','US/Eastern','US/Central');
+-----------------------------------------------------------+
```

```
| CONVERT_TZ('2007-03-11 3:00:00','US/Eastern','US/Central') |
+------------------------------------------------------------+
| 2007-03-11 01:00:00                                        |
+------------------------------------------------------------+
```

10.6.2 Time Zone Leap Second Support

Leap second values are returned with a time part that ends with `:59:59`. This means that a function such as `NOW()` can return the same value for two or three consecutive seconds during the leap second. It remains true that literal temporal values having a time part that ends with `:59:60` or `:59:61` are considered invalid.

If it is necessary to search for `TIMESTAMP` values one second before the leap second, anomalous results may be obtained if you use a comparison with `'YYYY-MM-DD hh:mm:ss'` values. The following example demonstrates this. It changes the local time zone to UTC so there is no difference between internal values (which are in UTC) and displayed values (which have time zone correction applied).

```
mysql> CREATE TABLE t1 (
         a INT,
         ts TIMESTAMP DEFAULT NOW(),
         PRIMARY KEY (ts)
       );
Query OK, 0 rows affected (0.01 sec)

mysql> -- change to UTC
mysql> SET time_zone = '+00:00';
Query OK, 0 rows affected (0.00 sec)

mysql> -- Simulate NOW() = '2008-12-31 23:59:59'
mysql> SET timestamp = 1230767999;
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t1 (a) VALUES (1);
Query OK, 1 row affected (0.00 sec)

mysql> -- Simulate NOW() = '2008-12-31 23:59:60'
mysql> SET timestamp = 1230768000;
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t1 (a) VALUES (2);
Query OK, 1 row affected (0.00 sec)

mysql> -- values differ internally but display the same
mysql> SELECT a, ts, UNIX_TIMESTAMP(ts) FROM t1;
+------+---------------------+--------------------+
| a    | ts                  | UNIX_TIMESTAMP(ts) |
+------+---------------------+--------------------+
|    1 | 2008-12-31 23:59:59 |         1230767999 |
|    2 | 2008-12-31 23:59:59 |         1230768000 |
+------+---------------------+--------------------+
2 rows in set (0.00 sec)

mysql> -- only the non-leap value matches
mysql> SELECT * FROM t1 WHERE ts = '2008-12-31 23:59:59';
+------+---------------------+
| a    | ts                  |
+------+---------------------+
|    1 | 2008-12-31 23:59:59 |
+------+---------------------+
1 row in set (0.00 sec)

mysql> -- the leap value with seconds=60 is invalid
mysql> SELECT * FROM t1 WHERE ts = '2008-12-31 23:59:60';
Empty set, 2 warnings (0.00 sec)
```

To work around this, you can use a comparison based on the UTC value actually stored in column, which has the leap second correction applied:

```
mysql> -- selecting using UNIX_TIMESTAMP value return leap value
mysql> SELECT * FROM t1 WHERE UNIX_TIMESTAMP(ts) = 1230768000;
+------+---------------------+
| a    | ts                  |
+------+---------------------+
|    2 | 2008-12-31 23:59:59 |
+------+---------------------+
1 row in set (0.00 sec)
```

10.7 MySQL Server Locale Support

The locale indicated by the `lc_time_names` system variable controls the language used to display day and month names and abbreviations. This variable affects the output from the `DATE_FORMAT()`, `DAYNAME()`, and `MONTHNAME()` functions.

`lc_time_names` does not affect the `STR_TO_DATE()` or `GET_FORMAT()` function.

The `lc_time_names` value does not affect the result from `FORMAT()`, but this function takes an optional third parameter that enables a locale to be specified to be used for the result number's decimal point, thousands separator, and grouping between separators. Permissible locale values are the same as the legal values for the `lc_time_names` system variable.

Locale names have language and region subtags listed by IANA (http://www.iana.org/assignments/language-subtag-registry) such as `'ja_JP'` or `'pt_BR'`. The default value is `'en_US'` regardless of your system's locale setting, but you can set the value at server startup or set the `GLOBAL` value if you have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege. Any client can examine the value of `lc_time_names` or set its `SESSION` value to affect the locale for its own connection.

```
mysql> SET NAMES 'utf8';
Query OK, 0 rows affected (0.09 sec)

mysql> SELECT @@lc_time_names;
+-----------------+
| @@lc_time_names |
+-----------------+
| en_US           |
+-----------------+
1 row in set (0.00 sec)

mysql> SELECT DAYNAME('2010-01-01'), MONTHNAME('2010-01-01');
+-----------------------+-------------------------+
| DAYNAME('2010-01-01') | MONTHNAME('2010-01-01') |
+-----------------------+-------------------------+
| Friday                | January                 |
+-----------------------+-------------------------+
1 row in set (0.00 sec)

mysql> SELECT DATE_FORMAT('2010-01-01','%W %a %M %b');
+-----------------------------------------+
| DATE_FORMAT('2010-01-01','%W %a %M %b') |
+-----------------------------------------+
| Friday Fri January Jan                  |
+-----------------------------------------+
1 row in set (0.00 sec)

mysql> SET lc_time_names = 'es_MX';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @@lc_time_names;
+-----------------+
| @@lc_time_names |
+-----------------+
| es_MX           |
+-----------------+
1 row in set (0.00 sec)
```

```
mysql> SELECT DAYNAME('2010-01-01'), MONTHNAME('2010-01-01');
+-----------------------+-------------------------+
| DAYNAME('2010-01-01') | MONTHNAME('2010-01-01') |
+-----------------------+-------------------------+
| viernes               | enero                   |
+-----------------------+-------------------------+
1 row in set (0.00 sec)

mysql> SELECT DATE_FORMAT('2010-01-01','%W %a %M %b');
+----------------------------------------+
| DATE_FORMAT('2010-01-01','%W %a %M %b') |
+----------------------------------------+
| viernes vie enero ene                  |
+----------------------------------------+
1 row in set (0.00 sec)
```

The day or month name for each of the affected functions is converted from `utf8` to the character set indicated by the `character_set_connection` system variable.

`lc_time_names` may be set to any of the following locale values. The set of locales supported by MySQL may differ from those supported by your operating system.

`ar_AE`: Arabic - United Arab Emirates	`ar_BH`: Arabic - Bahrain
`ar_DZ`: Arabic - Algeria	`ar_EG`: Arabic - Egypt
`ar_IN`: Arabic - India	`ar_IQ`: Arabic - Iraq
`ar_JO`: Arabic - Jordan	`ar_KW`: Arabic - Kuwait
`ar_LB`: Arabic - Lebanon	`ar_LY`: Arabic - Libya
`ar_MA`: Arabic - Morocco	`ar_OM`: Arabic - Oman
`ar_QA`: Arabic - Qatar	`ar_SA`: Arabic - Saudi Arabia
`ar_SD`: Arabic - Sudan	`ar_SY`: Arabic - Syria
`ar_TN`: Arabic - Tunisia	`ar_YE`: Arabic - Yemen
`be_BY`: Belarusian - Belarus	`bg_BG`: Bulgarian - Bulgaria
`ca_ES`: Catalan - Spain	`cs_CZ`: Czech - Czech Republic
`da_DK`: Danish - Denmark	`de_AT`: German - Austria
`de_BE`: German - Belgium	`de_CH`: German - Switzerland
`de_DE`: German - Germany	`de_LU`: German - Luxembourg
`el_GR`: Greek - Greece	`en_AU`: English - Australia
`en_CA`: English - Canada	`en_GB`: English - United Kingdom
`en_IN`: English - India	`en_NZ`: English - New Zealand
`en_PH`: English - Philippines	`en_US`: English - United States
`en_ZA`: English - South Africa	`en_ZW`: English - Zimbabwe
`es_AR`: Spanish - Argentina	`es_BO`: Spanish - Bolivia
`es_CL`: Spanish - Chile	`es_CO`: Spanish - Columbia
`es_CR`: Spanish - Costa Rica	`es_DO`: Spanish - Dominican Republic
`es_EC`: Spanish - Ecuador	`es_ES`: Spanish - Spain
`es_GT`: Spanish - Guatemala	`es_HN`: Spanish - Honduras
`es_MX`: Spanish - Mexico	`es_NI`: Spanish - Nicaragua
`es_PA`: Spanish - Panama	`es_PE`: Spanish - Peru
`es_PR`: Spanish - Puerto Rico	`es_PY`: Spanish - Paraguay
`es_SV`: Spanish - El Salvador	`es_US`: Spanish - United States
`es_UY`: Spanish - Uruguay	`es_VE`: Spanish - Venezuela

et_EE: Estonian - Estonia	eu_ES: Basque - Basque
fi_FI: Finnish - Finland	fo_FO: Faroese - Faroe Islands
fr_BE: French - Belgium	fr_CA: French - Canada
fr_CH: French - Switzerland	fr_FR: French - France
fr_LU: French - Luxembourg	gl_ES: Galician - Spain
gu_IN: Gujarati - India	he_IL: Hebrew - Israel
hi_IN: Hindi - India	hr_HR: Croatian - Croatia
hu_HU: Hungarian - Hungary	id_ID: Indonesian - Indonesia
is_IS: Icelandic - Iceland	it_CH: Italian - Switzerland
it_IT: Italian - Italy	ja_JP: Japanese - Japan
ko_KR: Korean - Republic of Korea	lt_LT: Lithuanian - Lithuania
lv_LV: Latvian - Latvia	mk_MK: Macedonian - FYROM
mn_MN: Mongolia - Mongolian	ms_MY: Malay - Malaysia
nb_NO: Norwegian(Bokmål) - Norway	nl_BE: Dutch - Belgium
nl_NL: Dutch - The Netherlands	no_NO: Norwegian - Norway
pl_PL: Polish - Poland	pt_BR: Portugese - Brazil
pt_PT: Portugese - Portugal	rm_CH: Romansh - Switzerland
ro_RO: Romanian - Romania	ru_RU: Russian - Russia
ru_UA: Russian - Ukraine	sk_SK: Slovak - Slovakia
sl_SI: Slovenian - Slovenia	sq_AL: Albanian - Albania
sr_RS: Serbian - Yugoslavia	sv_FI: Swedish - Finland
sv_SE: Swedish - Sweden	ta_IN: Tamil - India
te_IN: Telugu - India	th_TH: Thai - Thailand
tr_TR: Turkish - Turkey	uk_UA: Ukrainian - Ukraine
ur_PK: Urdu - Pakistan	vi_VN: Vietnamese - Viet Nam
zh_CN: Chinese - China	zh_HK: Chinese - Hong Kong
zh_TW: Chinese - Taiwan Province of China	

Chapter 11 Data Types

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MySQL supports a number of SQL data types in several categories: numeric types, date and time types, string (character and byte) types, spatial types, and the `JSON` data type. This chapter provides an overview of these data types, a more detailed description of the properties of the types in each category, and a summary of the data type storage requirements. The initial overview is intentionally brief. The more detailed descriptions later in the chapter should be consulted for additional information about particular data types, such as the permissible formats in which you can specify values.

Data type descriptions use these conventions:

- *M* indicates the maximum display width for integer types. For floating-point and fixed-point types, *M* is the total number of digits that can be stored (the precision). For string types, *M* is the maximum length. The maximum permissible value of *M* depends on the data type.
- *D* applies to floating-point and fixed-point types and indicates the number of digits following the decimal point (the scale). The maximum possible value is 30, but should be no greater than *M*−2.
- *fsp* applies to the `TIME`, `DATETIME`, and `TIMESTAMP` types and represents fractional seconds precision; that is, the number of digits following the decimal point for fractional parts of seconds. The *fsp* value, if given, must be in the range 0 to 6. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0. (This differs from the standard SQL default of 6, for compatibility with previous MySQL versions.)
- Square brackets (`[` and `]`) indicate optional parts of type definitions.

11.1 Data Type Overview

11.1.1 Numeric Type Overview

A summary of the numeric data types follows. For additional information about properties and storage requirements of the numeric types, see Section 11.2, “Numeric Types”, and Section 11.8, “Data Type Storage Requirements”.

M indicates the maximum display width for integer types. The maximum display width is 255. Display width is unrelated to the range of values a type can contain, as described in Section 11.2, “Numeric Types”. For floating-point and fixed-point types, *M* is the total number of digits that can be stored.

If you specify `ZEROFILL` for a numeric column, MySQL automatically adds the `UNSIGNED` attribute to the column.

Numeric data types that permit the `UNSIGNED` attribute also permit `SIGNED`. However, these data types are signed by default, so the `SIGNED` attribute has no effect.

`SERIAL` is an alias for `BIGINT UNSIGNED NOT NULL AUTO_INCREMENT UNIQUE`.

`SERIAL DEFAULT VALUE` in the definition of an integer column is an alias for `NOT NULL AUTO_INCREMENT UNIQUE`.

Warning

When you use subtraction between integer values where one is of type `UNSIGNED`, the result is unsigned unless the `NO_UNSIGNED_SUBTRACTION` SQL mode is enabled. See Section 12.10, “Cast Functions and Operators”.

- `BIT[(M)]`

  A bit-value type. *M* indicates the number of bits per value, from 1 to 64. The default is 1 if *M* is omitted.

- `TINYINT[(M)] [UNSIGNED] [ZEROFILL]`

  A very small integer. The signed range is `-128` to `127`. The unsigned range is `0` to `255`.

- `BOOL`, `BOOLEAN`

  These types are synonyms for `TINYINT(1)`. A value of zero is considered false. Nonzero values are considered true:

```
mysql> SELECT IF(0, 'true', 'false');
+-----------------------+
| IF(0, 'true', 'false') |
+-----------------------+
```

```
| false                  |
+------------------------+

mysql> SELECT IF(1, 'true', 'false');
+-----------------------+
| IF(1, 'true', 'false') |
+-----------------------+
| true                  |
+-----------------------+

mysql> SELECT IF(2, 'true', 'false');
+-----------------------+
| IF(2, 'true', 'false') |
+-----------------------+
| true                  |
+-----------------------+
```

However, the values `TRUE` and `FALSE` are merely aliases for `1` and `0`, respectively, as shown here:

```
mysql> SELECT IF(0 = FALSE, 'true', 'false');
+--------------------------------+
| IF(0 = FALSE, 'true', 'false') |
+--------------------------------+
| true                           |
+--------------------------------+

mysql> SELECT IF(1 = TRUE, 'true', 'false');
+-------------------------------+
| IF(1 = TRUE, 'true', 'false') |
+-------------------------------+
| true                          |
+-------------------------------+

mysql> SELECT IF(2 = TRUE, 'true', 'false');
+-------------------------------+
| IF(2 = TRUE, 'true', 'false') |
+-------------------------------+
| false                         |
+-------------------------------+

mysql> SELECT IF(2 = FALSE, 'true', 'false');
+--------------------------------+
| IF(2 = FALSE, 'true', 'false') |
+--------------------------------+
| false                          |
+--------------------------------+
```

The last two statements display the results shown because `2` is equal to neither `1` nor `0`.

- `SMALLINT[(M)] [UNSIGNED] [ZEROFILL]`

  A small integer. The signed range is `-32768` to `32767`. The unsigned range is `0` to `65535`.

- `MEDIUMINT[(M)] [UNSIGNED] [ZEROFILL]`

  A medium-sized integer. The signed range is `-8388608` to `8388607`. The unsigned range is `0` to `16777215`.

- `INT[(M)] [UNSIGNED] [ZEROFILL]`

  A normal-size integer. The signed range is `-2147483648` to `2147483647`. The unsigned range is `0` to `4294967295`.

- `INTEGER[(M)] [UNSIGNED] [ZEROFILL]`

  This type is a synonym for `INT`.

- `BIGINT[(M)] [UNSIGNED] [ZEROFILL]`

A large integer. The signed range is `-9223372036854775808` to `9223372036854775807`. The unsigned range is `0` to `18446744073709551615`.

`SERIAL` is an alias for `BIGINT UNSIGNED NOT NULL AUTO_INCREMENT UNIQUE`.

Some things you should be aware of with respect to `BIGINT` columns:

- All arithmetic is done using signed `BIGINT` or `DOUBLE` values, so you should not use unsigned big integers larger than `9223372036854775807` (63 bits) except with bit functions! If you do that, some of the last digits in the result may be wrong because of rounding errors when converting a `BIGINT` value to a `DOUBLE`.

  MySQL can handle `BIGINT` in the following cases:

  - When using integers to store large unsigned values in a `BIGINT` column.
  - In `MIN(col_name)` or `MAX(col_name)`, where `col_name` refers to a `BIGINT` column.
  - When using operators (`+`, `-`, `*`, and so on) where both operands are integers.

- You can always store an exact integer value in a `BIGINT` column by storing it using a string. In this case, MySQL performs a string-to-number conversion that involves no intermediate double-precision representation.
- The `-`, `+`, and `*` operators use `BIGINT` arithmetic when both operands are integer values. This means that if you multiply two big integers (or results from functions that return integers), you may get unexpected results when the result is larger than `9223372036854775807`.

- `DECIMAL[(M[,D])] [UNSIGNED] [ZEROFILL]`

  A packed “exact” fixed-point number. `M` is the total number of digits (the precision) and `D` is the number of digits after the decimal point (the scale). The decimal point and (for negative numbers) the `-` sign are not counted in `M`. If `D` is 0, values have no decimal point or fractional part. The maximum number of digits (`M`) for `DECIMAL` is 65. The maximum number of supported decimals (`D`) is 30. If `D` is omitted, the default is 0. If `M` is omitted, the default is 10.

  `UNSIGNED`, if specified, disallows negative values.

  All basic calculations (`+, -, *, /`) with `DECIMAL` columns are done with a precision of 65 digits.

- `DEC[(M[,D])] [UNSIGNED] [ZEROFILL]`, `NUMERIC[(M[,D])] [UNSIGNED] [ZEROFILL]`, `FIXED[(M[,D])] [UNSIGNED] [ZEROFILL]`

  These types are synonyms for `DECIMAL`. The `FIXED` synonym is available for compatibility with other database systems.

- `FLOAT[(M,D)] [UNSIGNED] [ZEROFILL]`

  A small (single-precision) floating-point number. Permissible values are `-3.402823466E+38` to `-1.175494351E-38`, `0`, and `1.175494351E-38` to `3.402823466E+38`. These are the theoretical limits, based on the IEEE standard. The actual range might be slightly smaller depending on your hardware or operating system.

  `M` is the total number of digits and `D` is the number of digits following the decimal point. If `M` and `D` are omitted, values are stored to the limits permitted by the hardware. A single-precision floating-point number is accurate to approximately 7 decimal places.

  `UNSIGNED`, if specified, disallows negative values.

  Using `FLOAT` might give you some unexpected problems because all calculations in MySQL are done with double precision. See Section B.5.4.7, “Solving Problems with No Matching Rows”.

- `DOUBLE[(M,D)] [UNSIGNED] [ZEROFILL]`

  A normal-size (double-precision) floating-point number. Permissible values are `-1.7976931348623157E+308` to `-2.2250738585072014E-308`, `0`, and `2.2250738585072014E-308` to `1.7976931348623157E+308`. These are the theoretical limits, based on the IEEE standard. The actual range might be slightly smaller depending on your hardware or operating system.

  M is the total number of digits and *D* is the number of digits following the decimal point. If *M* and *D* are omitted, values are stored to the limits permitted by the hardware. A double-precision floating-point number is accurate to approximately 15 decimal places.

  `UNSIGNED`, if specified, disallows negative values.

- `DOUBLE PRECISION[(M,D)] [UNSIGNED] [ZEROFILL]`, `REAL[(M,D)] [UNSIGNED] [ZEROFILL]`

  These types are synonyms for `DOUBLE`. Exception: If the `REAL_AS_FLOAT` SQL mode is enabled, `REAL` is a synonym for `FLOAT` rather than `DOUBLE`.

- `FLOAT(p) [UNSIGNED] [ZEROFILL]`

  A floating-point number. *p* represents the precision in bits, but MySQL uses this value only to determine whether to use `FLOAT` or `DOUBLE` for the resulting data type. If *p* is from 0 to 24, the data type becomes `FLOAT` with no *M* or *D* values. If *p* is from 25 to 53, the data type becomes `DOUBLE` with no *M* or *D* values. The range of the resulting column is the same as for the single-precision `FLOAT` or double-precision `DOUBLE` data types described earlier in this section.

  `FLOAT(p)` syntax is provided for ODBC compatibility.

11.1.2 Date and Time Type Overview

A summary of the temporal data types follows. For additional information about properties and storage requirements of the temporal types, see Section 11.3, “Date and Time Types”, and Section 11.8, “Data Type Storage Requirements”. For descriptions of functions that operate on temporal values, see Section 12.7, “Date and Time Functions”.

For the `DATE` and `DATETIME` range descriptions, “supported” means that although earlier values might work, there is no guarantee.

MySQL permits fractional seconds for `TIME`, `DATETIME`, and `TIMESTAMP` values, with up to microseconds (6 digits) precision. To define a column that includes a fractional seconds part, use the syntax *type_name(fsp)*, where *type_name* is `TIME`, `DATETIME`, or `TIMESTAMP`, and *fsp* is the fractional seconds precision. For example:

```
CREATE TABLE t1 (t TIME(3), dt DATETIME(6));
```

The *fsp* value, if given, must be in the range 0 to 6. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0. (This differs from the standard SQL default of 6, for compatibility with previous MySQL versions.)

Any `TIMESTAMP` or `DATETIME` column in a table can have automatic initialization and updating properties.

- `DATE`

  A date. The supported range is `'1000-01-01'` to `'9999-12-31'`. MySQL displays `DATE` values in `'YYYY-MM-DD'` format, but permits assignment of values to `DATE` columns using either strings or numbers.

- `DATETIME[(fsp)]`

A date and time combination. The supported range is `'1000-01-01 00:00:00.000000'` to `'9999-12-31 23:59:59.999999'`. MySQL displays `DATETIME` values in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format, but permits assignment of values to `DATETIME` columns using either strings or numbers.

An optional *`fsp`* value in the range from 0 to 6 may be given to specify fractional seconds precision. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0.

Automatic initialization and updating to the current date and time for `DATETIME` columns can be specified using `DEFAULT` and `ON UPDATE` column definition clauses, as described in Section 11.3.5, “Automatic Initialization and Updating for TIMESTAMP and DATETIME”.

- `TIMESTAMP[(`*`fsp`*`)]`

  A timestamp. The range is `'1970-01-01 00:00:01.000000'` UTC to `'2038-01-19 03:14:07.999999'` UTC. `TIMESTAMP` values are stored as the number of seconds since the epoch (`'1970-01-01 00:00:00'` UTC). A `TIMESTAMP` cannot represent the value `'1970-01-01 00:00:00'` because that is equivalent to 0 seconds from the epoch and the value 0 is reserved for representing `'0000-00-00 00:00:00'`, the “zero” `TIMESTAMP` value.

  An optional *`fsp`* value in the range from 0 to 6 may be given to specify fractional seconds precision. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0.

  The way the server handles `TIMESTAMP` definitions depends on the value of the `explicit_defaults_for_timestamp` system variable (see Section 5.1.5, “Server System Variables”).

  If `explicit_defaults_for_timestamp` is enabled, there is no automatic assignment of the `DEFAULT CURRENT_TIMESTAMP` or `ON UPDATE CURRENT_TIMESTAMP` attributes to any `TIMESTAMP` column. They must be included explicitly in the column definition. Also, any `TIMESTAMP` not explicitly declared as `NOT NULL` permits `NULL` values.

  If `explicit_defaults_for_timestamp` is disabled, the server handles `TIMESTAMP` as follows:

  Unless specified otherwise, the first `TIMESTAMP` column in a table is defined to be automatically set to the date and time of the most recent modification if not explicitly assigned a value. This makes `TIMESTAMP` useful for recording the timestamp of an `INSERT` or `UPDATE` operation. You can also set any `TIMESTAMP` column to the current date and time by assigning it a `NULL` value, unless it has been defined with the `NULL` attribute to permit `NULL` values.

  Automatic initialization and updating to the current date and time can be specified using `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP` column definition clauses. By default, the first `TIMESTAMP` column has these properties, as previously noted. However, any `TIMESTAMP` column in a table can be defined to have these properties.

- `TIME[(`*`fsp`*`)]`

  A time. The range is `'-838:59:59.000000'` to `'838:59:59.000000'`. MySQL displays `TIME` values in `'HH:MM:SS[.fraction]'` format, but permits assignment of values to `TIME` columns using either strings or numbers.

  An optional *`fsp`* value in the range from 0 to 6 may be given to specify fractional seconds precision. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0.

- `YEAR[(4)]`

  A year in four-digit format. MySQL displays `YEAR` values in `YYYY` format, but permits assignment of values to `YEAR` columns using either strings or numbers. Values display as `1901` to `2155`, and `0000`.

For additional information about YEAR display format and interpretation of input values, see Section 11.3.3, “The YEAR Type”.

Note

MySQL 8.0 does not support the YEAR(2) data type permitted in older versions of MySQL. For instructions on converting to YEAR(4), see Section 11.3.4, “Migrating YEAR(2) Columns to YEAR(4)”.

The SUM() and AVG() aggregate functions do not work with temporal values. (They convert the values to numbers, losing everything after the first nonnumeric character.) To work around this problem, convert to numeric units, perform the aggregate operation, and convert back to a temporal value. Examples:

```
SELECT SEC_TO_TIME(SUM(TIME_TO_SEC(time_col))) FROM tbl_name;
SELECT FROM_DAYS(SUM(TO_DAYS(date_col))) FROM tbl_name;
```

Note

The MySQL server can be run with the MAXDB SQL mode enabled. In this case, TIMESTAMP is identical with DATETIME. If this mode is enabled at the time that a table is created, TIMESTAMP columns are created as DATETIME columns. As a result, such columns use DATETIME display format, have the same range of values, and there is no automatic initialization or updating to the current date and time. See Section 5.1.8, “Server SQL Modes”.

11.1.3 String Type Overview

A summary of the string data types follows. For additional information about properties and storage requirements of the string types, see Section 11.4, “String Types”, and Section 11.8, “Data Type Storage Requirements”.

In some cases, MySQL may change a string column to a type different from that given in a CREATE TABLE or ALTER TABLE statement. See Section 13.1.15.7, “Silent Column Specification Changes”.

MySQL interprets length specifications in character column definitions in character units. This applies to CHAR, VARCHAR, and the TEXT types.

Column definitions for many string data types can include attributes that specify the character set or collation of the column. These attributes apply to the CHAR, VARCHAR, the TEXT types, ENUM, and SET data types:

- The CHARACTER SET attribute specifies the character set, and the COLLATE attribute specifies a collation for the character set. For example:

  ```
  CREATE TABLE t
  (
      c1 VARCHAR(20) CHARACTER SET utf8,
      c2 TEXT CHARACTER SET latin1 COLLATE latin1_general_cs
  );
  ```

  This table definition creates a column named c1 that has a character set of utf8 with the default collation for that character set, and a column named c2 that has a character set of latin1 and a case-sensitive collation.

  The rules for assigning the character set and collation when either or both of the CHARACTER SET and COLLATE attributes are missing are described in Section 10.1.3.5, “Column Character Set and Collation”.

  CHARSET is a synonym for CHARACTER SET.

- Specifying the `CHARACTER SET binary` attribute for a character string data type causes the column to be created as the corresponding binary string data type: `CHAR` becomes `BINARY`, `VARCHAR` becomes `VARBINARY`, and `TEXT` becomes `BLOB`. For the `ENUM` and `SET` data types, this does not occur; they are created as declared. Suppose that you specify a table using this definition:

```
CREATE TABLE t
(
  c1 VARCHAR(10) CHARACTER SET binary,
  c2 TEXT CHARACTER SET binary,
  c3 ENUM('a','b','c') CHARACTER SET binary
);
```

  The resulting table has this definition:

```
CREATE TABLE t
(
  c1 VARBINARY(10),
  c2 BLOB,
  c3 ENUM('a','b','c') CHARACTER SET binary
);
```

- The `BINARY` attribute is shorthand for specifying the table default character set and the binary (`_bin`) collation of that character set. In this case, comparison and sorting are based on numeric character code values.
- The `ASCII` attribute is shorthand for `CHARACTER SET latin1`.
- The `UNICODE` attribute is shorthand for `CHARACTER SET ucs2`.

Character column comparison and sorting are based on the collation assigned to the column. For the `CHAR`, `VARCHAR`, `TEXT`, `ENUM`, and `SET` data types, you can declare a column with a binary (`_bin`) collation or the `BINARY` attribute to cause comparison and sorting to use the underlying character code values rather than a lexical ordering.

For additional information about use of character sets in MySQL, see Section 10.1, “Character Set Support”.

- `[NATIONAL] CHAR[(M)] [CHARACTER SET charset_name] [COLLATE collation_name]`

  A fixed-length string that is always right-padded with spaces to the specified length when stored. *`M`* represents the column length in characters. The range of *`M`* is 0 to 255. If *`M`* is omitted, the length is 1.

  Note

  Trailing spaces are removed when `CHAR` values are retrieved unless the `PAD_CHAR_TO_FULL_LENGTH` SQL mode is enabled.

  `CHAR` is shorthand for `CHARACTER`. `NATIONAL CHAR` (or its equivalent short form, `NCHAR`) is the standard SQL way to define that a `CHAR` column should use some predefined character set. MySQL uses `utf8` as this predefined character set. Section 10.1.3.7, “The National Character Set”.

  The `CHAR BYTE` data type is an alias for the `BINARY` data type. This is a compatibility feature.

  MySQL permits you to create a column of type `CHAR(0)`. This is useful primarily when you have to be compliant with old applications that depend on the existence of a column but that do not actually use its value. `CHAR(0)` is also quite nice when you need a column that can take only two values: A column that is defined as `CHAR(0) NULL` occupies only one bit and can take only the values `NULL` and `''` (the empty string).

- `[NATIONAL] VARCHAR(M) [CHARACTER SET charset_name] [COLLATE collation_name]`

A variable-length string. *M* represents the maximum column length in characters. The range of *M* is 0 to 65,535. The effective maximum length of a VARCHAR is subject to the maximum row size (65,535 bytes, which is shared among all columns) and the character set used. For example, utf8 characters can require up to three bytes per character, so a VARCHAR column that uses the utf8 character set can be declared to be a maximum of 21,844 characters. See Section C.10.4, "Limits on Table Column Count and Row Size".

MySQL stores VARCHAR values as a 1-byte or 2-byte length prefix plus data. The length prefix indicates the number of bytes in the value. A VARCHAR column uses one length byte if values require no more than 255 bytes, two length bytes if values may require more than 255 bytes.

Note

MySQL follows the standard SQL specification, and does *not* remove trailing spaces from VARCHAR values.

VARCHAR is shorthand for CHARACTER VARYING. NATIONAL VARCHAR is the standard SQL way to define that a VARCHAR column should use some predefined character set. MySQL uses utf8 as this predefined character set. Section 10.1.3.7, "The National Character Set". NVARCHAR is shorthand for NATIONAL VARCHAR.

- BINARY(*M*)

  The BINARY type is similar to the CHAR type, but stores binary byte strings rather than nonbinary character strings. *M* represents the column length in bytes.

- VARBINARY(*M*)

  The VARBINARY type is similar to the VARCHAR type, but stores binary byte strings rather than nonbinary character strings. *M* represents the maximum column length in bytes.

- TINYBLOB

  A BLOB column with a maximum length of 255 (2^8 − 1) bytes. Each TINYBLOB value is stored using a 1-byte length prefix that indicates the number of bytes in the value.

- TINYTEXT [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  A TEXT column with a maximum length of 255 (2^8 − 1) characters. The effective maximum length is less if the value contains multibyte characters. Each TINYTEXT value is stored using a 1-byte length prefix that indicates the number of bytes in the value.

- BLOB[(*M*)]

  A BLOB column with a maximum length of 65,535 (2^{16} − 1) bytes. Each BLOB value is stored using a 2-byte length prefix that indicates the number of bytes in the value.

  An optional length *M* can be given for this type. If this is done, MySQL creates the column as the smallest BLOB type large enough to hold values *M* bytes long.

- TEXT[(*M*)] [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  A TEXT column with a maximum length of 65,535 (2^{16} − 1) characters. The effective maximum length is less if the value contains multibyte characters. Each TEXT value is stored using a 2-byte length prefix that indicates the number of bytes in the value.

  An optional length *M* can be given for this type. If this is done, MySQL creates the column as the smallest TEXT type large enough to hold values *M* characters long.

- MEDIUMBLOB

A BLOB column with a maximum length of 16,777,215 (2^{24} − 1) bytes. Each MEDIUMBLOB value is stored using a 3-byte length prefix that indicates the number of bytes in the value.

- MEDIUMTEXT [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  A TEXT column with a maximum length of 16,777,215 (2^{24} − 1) characters. The effective maximum length is less if the value contains multibyte characters. Each MEDIUMTEXT value is stored using a 3-byte length prefix that indicates the number of bytes in the value.

- LONGBLOB

  A BLOB column with a maximum length of 4,294,967,295 or 4GB (2^{32} − 1) bytes. The effective maximum length of LONGBLOB columns depends on the configured maximum packet size in the client/server protocol and available memory. Each LONGBLOB value is stored using a 4-byte length prefix that indicates the number of bytes in the value.

- LONGTEXT [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  A TEXT column with a maximum length of 4,294,967,295 or 4GB (2^{32} − 1) characters. The effective maximum length is less if the value contains multibyte characters. The effective maximum length of LONGTEXT columns also depends on the configured maximum packet size in the client/server protocol and available memory. Each LONGTEXT value is stored using a 4-byte length prefix that indicates the number of bytes in the value.

- ENUM('*value1*','*value2*',...) [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  An enumeration. A string object that can have only one value, chosen from the list of values '*value1*', '*value2*', ..., NULL or the special '' error value. ENUM values are represented internally as integers.

  An ENUM column can have a maximum of 65,535 distinct elements.

  As of MySQL 8.0.2, the maximum supported length of an individual ENUM element is *M* <= 255 and (*M* x *w*) <= 1020, where M is the element literal length and *w* is the number of bytes required for the maximum-length character in the character set.

- SET('*value1*','*value2*',...) [CHARACTER SET *charset_name*] [COLLATE *collation_name*]

  A set. A string object that can have zero or more values, each of which must be chosen from the list of values '*value1*', '*value2*', ... SET values are represented internally as integers.

  A SET column can have a maximum of 64 distinct members.

  As of MySQL 8.0.2, the maximum supported length of an individual SET element is *M* <= 255 and (*M* x *w*) <= 1020, where M is the element literal length and *w* is the number of bytes required for the maximum-length character in the character set.

11.2 Numeric Types

MySQL supports all standard SQL numeric data types. These types include the exact numeric data types (INTEGER, SMALLINT, DECIMAL, and NUMERIC), as well as the approximate numeric data types (FLOAT, REAL, and DOUBLE PRECISION). The keyword INT is a synonym for INTEGER, and the keywords DEC and FIXED are synonyms for DECIMAL. MySQL treats DOUBLE as a synonym for DOUBLE PRECISION (a nonstandard extension). MySQL also treats REAL as a synonym for DOUBLE PRECISION (a nonstandard variation), unless the REAL_AS_FLOAT SQL mode is enabled.

The BIT data type stores bit values and is supported for MyISAM, MEMORY, and InnoDB.

For information about how MySQL handles assignment of out-of-range values to columns and overflow during expression evaluation, see Section 11.2.6, “Out-of-Range and Overflow Handling”.

For information about numeric type storage requirements, see Section 11.8, “Data Type Storage Requirements”.

The data type used for the result of a calculation on numeric operands depends on the types of the operands and the operations performed on them. For more information, see Section 12.6.1, “Arithmetic Operators”.

11.2.1 Integer Types (Exact Value) - INTEGER, INT, SMALLINT, TINYINT, MEDIUMINT, BIGINT

MySQL supports the SQL standard integer types `INTEGER` (or `INT`) and `SMALLINT`. As an extension to the standard, MySQL also supports the integer types `TINYINT`, `MEDIUMINT`, and `BIGINT`. The following table shows the required storage and range for each integer type.

Type	Storage	Minimum Value	Maximum Value
	(Bytes)	(Signed/Unsigned)	(Signed/Unsigned)
`TINYINT`	1	`-128`	`127`
		`0`	`255`
`SMALLINT`	2	`-32768`	`32767`
		`0`	`65535`
`MEDIUMINT`	3	`-8388608`	`8388607`
		`0`	`16777215`
`INT`	4	`-2147483648`	`2147483647`
		`0`	`4294967295`
`BIGINT`	8	`-9223372036854775808`	`9223372036854775807`
		`0`	`18446744073709551615`

11.2.2 Fixed-Point Types (Exact Value) - DECIMAL, NUMERIC

The `DECIMAL` and `NUMERIC` types store exact numeric data values. These types are used when it is important to preserve exact precision, for example with monetary data. In MySQL, `NUMERIC` is implemented as `DECIMAL`, so the following remarks about `DECIMAL` apply equally to `NUMERIC`.

MySQL stores `DECIMAL` values in binary format. See Section 12.21, “Precision Math”.

In a `DECIMAL` column declaration, the precision and scale can be (and usually is) specified; for example:

```
salary DECIMAL(5,2)
```

In this example, `5` is the precision and `2` is the scale. The precision represents the number of significant digits that are stored for values, and the scale represents the number of digits that can be stored following the decimal point.

Standard SQL requires that `DECIMAL(5,2)` be able to store any value with five digits and two decimals, so values that can be stored in the `salary` column range from `-999.99` to `999.99`.

In standard SQL, the syntax `DECIMAL(M)` is equivalent to `DECIMAL(M,0)`. Similarly, the syntax `DECIMAL` is equivalent to `DECIMAL(M,0)`, where the implementation is permitted to decide the value of *`M`*. MySQL supports both of these variant forms of `DECIMAL` syntax. The default value of *`M`* is 10.

If the scale is 0, `DECIMAL` values contain no decimal point or fractional part.

The maximum number of digits for `DECIMAL` is 65, but the actual range for a given `DECIMAL` column can be constrained by the precision or scale for a given column. When such a column is assigned a value with more digits following the decimal point than are permitted by the specified scale, the value is converted to that scale. (The precise behavior is operating system-specific, but generally the effect is truncation to the permissible number of digits.)

11.2.3 Floating-Point Types (Approximate Value) - FLOAT, DOUBLE

The `FLOAT` and `DOUBLE` types represent approximate numeric data values. MySQL uses four bytes for single-precision values and eight bytes for double-precision values.

For `FLOAT`, the SQL standard permits an optional specification of the precision (but not the range of the exponent) in bits following the keyword `FLOAT` in parentheses. MySQL also supports this optional precision specification, but the precision value is used only to determine storage size. A precision from 0 to 23 results in a 4-byte single-precision `FLOAT` column. A precision from 24 to 53 results in an 8-byte double-precision `DOUBLE` column.

MySQL permits a nonstandard syntax: `FLOAT(M,D)` or `REAL(M,D)` or `DOUBLE PRECISION(M,D)`. Here, `(M,D)` means than values can be stored with up to *`M`* digits in total, of which *`D`* digits may be after the decimal point. For example, a column defined as `FLOAT(7,4)` will look like `-999.9999` when displayed. MySQL performs rounding when storing values, so if you insert `999.00009` into a `FLOAT(7,4)` column, the approximate result is `999.0001`.

Because floating-point values are approximate and not stored as exact values, attempts to treat them as exact in comparisons may lead to problems. They are also subject to platform or implementation dependencies. For more information, see Section B.5.4.8, “Problems with Floating-Point Values”

For maximum portability, code requiring storage of approximate numeric data values should use `FLOAT` or `DOUBLE PRECISION` with no specification of precision or number of digits.

11.2.4 Bit-Value Type - BIT

The `BIT` data type is used to store bit values. A type of `BIT(M)` enables storage of *`M`*-bit values. *`M`* can range from 1 to 64.

To specify bit values, `b'value'` notation can be used. *`value`* is a binary value written using zeros and ones. For example, `b'111'` and `b'10000000'` represent 7 and 128, respectively. See Section 9.1.5, “Bit-Value Literals”.

If you assign a value to a `BIT(M)` column that is less than *`M`* bits long, the value is padded on the left with zeros. For example, assigning a value of `b'101'` to a `BIT(6)` column is, in effect, the same as assigning `b'000101'`.

11.2.5 Numeric Type Attributes

MySQL supports an extension for optionally specifying the display width of integer data types in parentheses following the base keyword for the type. For example, `INT(4)` specifies an `INT` with a display width of four digits. This optional display width may be used by applications to display integer values having a width less than the width specified for the column by left-padding them with spaces. (That is, this width is present in the metadata returned with result sets. Whether it is used or not is up to the application.)

The display width does *not* constrain the range of values that can be stored in the column. Nor does it prevent values wider than the column display width from being displayed correctly. For example, a column specified as `SMALLINT(3)` has the usual `SMALLINT` range of `-32768` to `32767`, and values outside the range permitted by three digits are displayed in full using more than three digits.

When used in conjunction with the optional (nonstandard) attribute `ZEROFILL`, the default padding of spaces is replaced with zeros. For example, for a column declared as `INT(4) ZEROFILL`, a value of `5` is retrieved as `0005`.

Note

The `ZEROFILL` attribute is ignored when a column is involved in expressions or `UNION` queries.

If you store values larger than the display width in an integer column that has the `ZEROFILL` attribute, you may experience problems when MySQL generates temporary tables for some complicated joins. In these cases, MySQL assumes that the data values fit within the column display width.

All integer types can have an optional (nonstandard) attribute `UNSIGNED`. Unsigned type can be used to permit only nonnegative numbers in a column or when you need a larger upper numeric range for the column. For example, if an `INT` column is `UNSIGNED`, the size of the column's range is the same but its endpoints shift from `-2147483648` and `2147483647` up to `0` and `4294967295`.

Floating-point and fixed-point types also can be `UNSIGNED`. As with integer types, this attribute prevents negative values from being stored in the column. Unlike the integer types, the upper range of column values remains the same.

If you specify `ZEROFILL` for a numeric column, MySQL automatically adds the `UNSIGNED` attribute to the column.

Integer or floating-point data types can have the additional attribute `AUTO_INCREMENT`. When you insert a value of `NULL` into an indexed `AUTO_INCREMENT` column, the column is set to the next sequence value. Typically this is *`value`*`+1`, where *`value`* is the largest value for the column currently in the table. (`AUTO_INCREMENT` sequences begin with `1`.)

Storing `0` into an `AUTO_INCREMENT` column has the same effect as storing `NULL`, unless the `NO_AUTO_VALUE_ON_ZERO` SQL mode is enabled.

Inserting `NULL` to generate `AUTO_INCREMENT` values requires that the column be declared `NOT NULL`. If the column is declared `NULL`, inserting `NULL` stores a `NULL`. When you insert any other value into an `AUTO_INCREMENT` column, the column is set to that value and the sequence is reset so that the next automatically generated value follows sequentially from the inserted value.

In MySQL 8.0, negative values for `AUTO_INCREMENT` columns are not supported.

11.2.6 Out-of-Range and Overflow Handling

When MySQL stores a value in a numeric column that is outside the permissible range of the column data type, the result depends on the SQL mode in effect at the time:

- If strict SQL mode is enabled, MySQL rejects the out-of-range value with an error, and the insert fails, in accordance with the SQL standard.
- If no restrictive modes are enabled, MySQL clips the value to the appropriate endpoint of the column data type range and stores the resulting value instead.

  When an out-of-range value is assigned to an integer column, MySQL stores the value representing the corresponding endpoint of the column data type range.

  When a floating-point or fixed-point column is assigned a value that exceeds the range implied by the specified (or default) precision and scale, MySQL stores the value representing the corresponding endpoint of that range.

Suppose that a table `t1` has this definition:

```
CREATE TABLE t1 (i1 TINYINT, i2 TINYINT UNSIGNED);
```

With strict SQL mode enabled, an out of range error occurs:

```
mysql> SET sql_mode = 'TRADITIONAL';
mysql> INSERT INTO t1 (i1, i2) VALUES(256, 256);
ERROR 1264 (22003): Out of range value for column 'i1' at row 1
mysql> SELECT * FROM t1;
Empty set (0.00 sec)
```

With strict SQL mode not enabled, clipping with warnings occurs:

```
mysql> SET sql_mode = '';
mysql> INSERT INTO t1 (i1, i2) VALUES(256, 256);
mysql> SHOW WARNINGS;
+---------+------+-----------------------------------------+
| Level   | Code | Message                                 |
+---------+------+-----------------------------------------+
| Warning | 1264 | Out of range value for column 'i1' at row 1 |
| Warning | 1264 | Out of range value for column 'i2' at row 1 |
+---------+------+-----------------------------------------+
mysql> SELECT * FROM t1;
+------+------+
| i1   | i2   |
+------+------+
|  127 |  255 |
+------+------+
```

When strict SQL mode is not enabled, column-assignment conversions that occur due to clipping are reported as warnings for `ALTER TABLE`, `LOAD DATA INFILE`, `UPDATE`, and multiple-row `INSERT` statements. In strict mode, these statements fail, and some or all the values are not inserted or changed, depending on whether the table is a transactional table and other factors. For details, see Section 5.1.8, “Server SQL Modes”.

Overflow during numeric expression evaluation results in an error. For example, the largest signed `BIGINT` value is 9223372036854775807, so the following expression produces an error:

```
mysql> SELECT 9223372036854775807 + 1;
ERROR 1690 (22003): BIGINT value is out of range in '(9223372036854775807 + 1)'
```

To enable the operation to succeed in this case, convert the value to unsigned;

```
mysql> SELECT CAST(9223372036854775807 AS UNSIGNED) + 1;
+-----------------------------------------+
| CAST(9223372036854775807 AS UNSIGNED) + 1 |
+-----------------------------------------+
|                      9223372036854775808 |
+-----------------------------------------+
```

Whether overflow occurs depends on the range of the operands, so another way to handle the preceding expression is to use exact-value arithmetic because `DECIMAL` values have a larger range than integers:

```
mysql> SELECT 9223372036854775807.0 + 1;
+--------------------------+
| 9223372036854775807.0 + 1 |
+--------------------------+
|     9223372036854775808.0 |
+--------------------------+
```

Subtraction between integer values, where one is of type `UNSIGNED`, produces an unsigned result by default. If the result would otherwise have been negative, an error results:

```
mysql> SET sql_mode = '';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT CAST(0 AS UNSIGNED) - 1;
ERROR 1690 (22003): BIGINT UNSIGNED value is out of range in '(cast(0 as unsigned) - 1)'
```

If the NO_UNSIGNED_SUBTRACTION SQL mode is enabled, the result is negative:

```
mysql> SET sql_mode = 'NO_UNSIGNED_SUBTRACTION';
mysql> SELECT CAST(0 AS UNSIGNED) - 1;
+--------------------------+
| CAST(0 AS UNSIGNED) - 1 |
+--------------------------+
|                       -1 |
+--------------------------+
```

If the result of such an operation is used to update an UNSIGNED integer column, the result is clipped to the maximum value for the column type, or clipped to 0 if NO_UNSIGNED_SUBTRACTION is enabled. If strict SQL mode is enabled, an error occurs and the column remains unchanged.

11.3 Date and Time Types

The date and time types for representing temporal values are DATE, TIME, DATETIME, TIMESTAMP, and YEAR. Each temporal type has a range of valid values, as well as a “zero” value that may be used when you specify an invalid value that MySQL cannot represent. The TIMESTAMP type has special automatic updating behavior, described later. For temporal type storage requirements, see Section 11.8, “Data Type Storage Requirements”.

Keep in mind these general considerations when working with date and time types:

- MySQL retrieves values for a given date or time type in a standard output format, but it attempts to interpret a variety of formats for input values that you supply (for example, when you specify a value to be assigned to or compared to a date or time type). For a description of the permitted formats for date and time types, see Section 9.1.3, “Date and Time Literals”. It is expected that you supply valid values. Unpredictable results may occur if you use values in other formats.

- Although MySQL tries to interpret values in several formats, date parts must always be given in year-month-day order (for example, '98-09-04'), rather than in the month-day-year or day-month-year orders commonly used elsewhere (for example, '09-04-98', '04-09-98').

- Dates containing two-digit year values are ambiguous because the century is unknown. MySQL interprets two-digit year values using these rules:

  - Year values in the range 70-99 are converted to 1970-1999.

  - Year values in the range 00-69 are converted to 2000-2069.

  See also Section 11.3.8, “Two-Digit Years in Dates”.

- Conversion of values from one temporal type to another occurs according to the rules in Section 11.3.7, “Conversion Between Date and Time Types”.

- MySQL automatically converts a date or time value to a number if the value is used in a numeric context and vice versa.

- By default, when MySQL encounters a value for a date or time type that is out of range or otherwise invalid for the type, it converts the value to the “zero” value for that type. The exception is that out-of-range TIME values are clipped to the appropriate endpoint of the TIME range.

- By setting the SQL mode to the appropriate value, you can specify more exactly what kind of dates you want MySQL to support. (See Section 5.1.8, “Server SQL Modes”.) You can get MySQL to

accept certain dates, such as `'2009-11-31'`, by enabling the `ALLOW_INVALID_DATES` SQL mode. This is useful when you want to store a “possibly wrong” value which the user has specified (for example, in a web form) in the database for future processing. Under this mode, MySQL verifies only that the month is in the range from 1 to 12 and that the day is in the range from 1 to 31.

- MySQL permits you to store dates where the day or month and day are zero in a `DATE` or `DATETIME` column. This is useful for applications that need to store birthdates for which you may not know the exact date. In this case, you simply store the date as `'2009-00-00'` or `'2009-01-00'`. If you store dates such as these, you should not expect to get correct results for functions such as `DATE_SUB()` or `DATE_ADD()` that require complete dates. To disallow zero month or day parts in dates, enable the `NO_ZERO_IN_DATE` mode.
- MySQL permits you to store a “zero” value of `'0000-00-00'` as a “dummy date.” This is in some cases more convenient than using `NULL` values, and uses less data and index space. To disallow `'0000-00-00'`, enable the `NO_ZERO_DATE` mode.
- “Zero” date or time values used through Connector/ODBC are converted automatically to `NULL` because ODBC cannot handle such values.

The following table shows the format of the “zero” value for each type. The “zero” values are special, but you can store or refer to them explicitly using the values shown in the table. You can also do this using the values `'0'` or `0`, which are easier to write. For temporal types that include a date part (`DATE`, `DATETIME`, and `TIMESTAMP`), use of these values produces warnings if the `NO_ZERO_DATE` SQL mode is enabled.

Data Type	“Zero” Value
`DATE`	`'0000-00-00'`
`TIME`	`'00:00:00'`
`DATETIME`	`'0000-00-00 00:00:00'`
`TIMESTAMP`	`'0000-00-00 00:00:00'`
`YEAR`	`0000`

11.3.1 The DATE, DATETIME, and TIMESTAMP Types

The `DATE`, `DATETIME`, and `TIMESTAMP` types are related. This section describes their characteristics, how they are similar, and how they differ. MySQL recognizes `DATE`, `DATETIME`, and `TIMESTAMP` values in several formats, described in Section 9.1.3, “Date and Time Literals”. For the `DATE` and `DATETIME` range descriptions, “supported” means that although earlier values might work, there is no guarantee.

The `DATE` type is used for values with a date part but no time part. MySQL retrieves and displays `DATE` values in `'YYYY-MM-DD'` format. The supported range is `'1000-01-01'` to `'9999-12-31'`.

The `DATETIME` type is used for values that contain both date and time parts. MySQL retrieves and displays `DATETIME` values in `'YYYY-MM-DD HH:MM:SS'` format. The supported range is `'1000-01-01 00:00:00'` to `'9999-12-31 23:59:59'`.

The `TIMESTAMP` data type is used for values that contain both date and time parts. `TIMESTAMP` has a range of `'1970-01-01 00:00:01'` UTC to `'2038-01-19 03:14:07'` UTC.

A `DATETIME` or `TIMESTAMP` value can include a trailing fractional seconds part in up to microseconds (6 digits) precision. In particular, any fractional part in a value inserted into a `DATETIME` or `TIMESTAMP` column is stored rather than discarded. With the fractional part included, the format for these values is `'YYYY-MM-DD HH:MM:SS[.fraction]'`, the range for `DATETIME` values is `'1000-01-01 00:00:00.000000'` to `'9999-12-31 23:59:59.999999'`, and the range for `TIMESTAMP` values is `'1970-01-01 00:00:01.000000'` to `'2038-01-19 03:14:07.999999'`. The fractional part should always be separated from the rest of the time by a decimal point; no other fractional

seconds delimiter is recognized. For information about fractional seconds support in MySQL, see Section 11.3.6, “Fractional Seconds in Time Values”.

The `TIMESTAMP` and `DATETIME` data types offer automatic initialization and updating to the current date and time. For more information, see Section 11.3.5, “Automatic Initialization and Updating for TIMESTAMP and DATETIME”.

MySQL converts `TIMESTAMP` values from the current time zone to UTC for storage, and back from UTC to the current time zone for retrieval. (This does not occur for other types such as `DATETIME`.) By default, the current time zone for each connection is the server's time. The time zone can be set on a per-connection basis. As long as the time zone setting remains constant, you get back the same value you store. If you store a `TIMESTAMP` value, and then change the time zone and retrieve the value, the retrieved value is different from the value you stored. This occurs because the same time zone was not used for conversion in both directions. The current time zone is available as the value of the `time_zone` system variable. For more information, see Section 10.6, “MySQL Server Time Zone Support”.

Invalid `DATE`, `DATETIME`, or `TIMESTAMP` values are converted to the “zero” value of the appropriate type (`'0000-00-00'` or `'0000-00-00 00:00:00'`).

Be aware of certain properties of date value interpretation in MySQL:

- MySQL permits a “relaxed” format for values specified as strings, in which any punctuation character may be used as the delimiter between date parts or time parts. In some cases, this syntax can be deceiving. For example, a value such as `'10:11:12'` might look like a time value because of the `:`, but is interpreted as the year `'2010-11-12'` if used in a date context. The value `'10:45:15'` is converted to `'0000-00-00'` because `'45'` is not a valid month.

  The only delimiter recognized between a date and time part and a fractional seconds part is the decimal point.

- The server requires that month and day values be valid, and not merely in the range 1 to 12 and 1 to 31, respectively. With strict mode disabled, invalid dates such as `'2004-04-31'` are converted to `'0000-00-00'` and a warning is generated. With strict mode enabled, invalid dates generate an error. To permit such dates, enable `ALLOW_INVALID_DATES`. See Section 5.1.8, “Server SQL Modes”, for more information.

- MySQL does not accept `TIMESTAMP` values that include a zero in the day or month column or values that are not a valid date. The sole exception to this rule is the special “zero” value `'0000-00-00 00:00:00'`.

- Dates containing two-digit year values are ambiguous because the century is unknown. MySQL interprets two-digit year values using these rules:

  - Year values in the range `00-69` are converted to `2000-2069`.
  - Year values in the range `70-99` are converted to `1970-1999`.

  See also Section 11.3.8, “Two-Digit Years in Dates”.

Note

The MySQL server can be run with the `MAXDB` SQL mode enabled. In this case, `TIMESTAMP` is identical with `DATETIME`. If this mode is enabled at the time that a table is created, `TIMESTAMP` columns are created as `DATETIME` columns. As a result, such columns use `DATETIME` display format, have the same range of values, and there is no automatic initialization or updating to the current date and time. See Section 5.1.8, “Server SQL Modes”.

11.3.2 The TIME Type

MySQL retrieves and displays `TIME` values in `'HH:MM:SS'` format (or `'HHH:MM:SS'` format for large hours values). `TIME` values may range from `'-838:59:59'` to `'838:59:59'`. The hours part may be so large because the `TIME` type can be used not only to represent a time of day (which must be less than 24 hours), but also elapsed time or a time interval between two events (which may be much greater than 24 hours, or even negative).

MySQL recognizes `TIME` values in several formats, some of which can include a trailing fractional seconds part in up to microseconds (6 digits) precision. See Section 9.1.3, “Date and Time Literals”. For information about fractional seconds support in MySQL, see Section 11.3.6, “Fractional Seconds in Time Values”. In particular, any fractional part in a value inserted into a `TIME` column is stored rather than discarded. With the fractional part included, the range for `TIME` values is `'-838:59:59.000000'` to `'838:59:59.000000'`.

Be careful about assigning abbreviated values to a `TIME` column. MySQL interprets abbreviated `TIME` values with colons as time of the day. That is, `'11:12'` means `'11:12:00'`, not `'00:11:12'`. MySQL interprets abbreviated values without colons using the assumption that the two rightmost digits represent seconds (that is, as elapsed time rather than as time of day). For example, you might think of `'1112'` and `1112` as meaning `'11:12:00'` (12 minutes after 11 o'clock), but MySQL interprets them as `'00:11:12'` (11 minutes, 12 seconds). Similarly, `'12'` and `12` are interpreted as `'00:00:12'`.

The only delimiter recognized between a time part and a fractional seconds part is the decimal point.

By default, values that lie outside the `TIME` range but are otherwise valid are clipped to the closest endpoint of the range. For example, `'-850:00:00'` and `'850:00:00'` are converted to `'-838:59:59'` and `'838:59:59'`. Invalid `TIME` values are converted to `'00:00:00'`. Note that because `'00:00:00'` is itself a valid `TIME` value, there is no way to tell, from a value of `'00:00:00'` stored in a table, whether the original value was specified as `'00:00:00'` or whether it was invalid.

For more restrictive treatment of invalid `TIME` values, enable strict SQL mode to cause errors to occur. See Section 5.1.8, “Server SQL Modes”.

11.3.3 The YEAR Type

The `YEAR` type is a 1-byte type used to represent year values. It can be declared as `YEAR` or `YEAR(4)` and has a display width of four characters.

Note

MySQL 8.0 does not support the `YEAR(2)` data type permitted in older versions of MySQL. For instructions on converting to `YEAR(4)`, see Section 11.3.4, “Migrating YEAR(2) Columns to YEAR(4)”.

MySQL displays `YEAR` values in `YYYY` format, with a range of `1901` to `2155`, or `0000`.

You can specify input `YEAR` values in a variety of formats:

- As a 4-digit number in the range `1901` to `2155`.
- As a 4-digit string in the range `'1901'` to `'2155'`.
- As a 1- or 2-digit number in the range `1` to `99`. MySQL converts values in the ranges `1` to `69` and `70` to `99` to `YEAR` values in the ranges `2001` to `2069` and `1970` to `1999`.
- As a 1- or 2-digit string in the range `'0'` to `'99'`. MySQL converts values in the ranges `'0'` to `'69'` and `'70'` to `'99'` to `YEAR` values in the ranges `2000` to `2069` and `1970` to `1999`.
- The result of inserting a numeric `0` has a display value of `0000` and an internal value of `0000`. To insert zero and have it be interpreted as `2000`, specify it as a string `'0'` or `'00'`.
- As the result of a function that returns a value that is acceptable in a `YEAR` context, such as `NOW()`.

MySQL converts invalid `YEAR` values to `0000`.

See also Section 11.3.8, “Two-Digit Years in Dates”.

11.3.4 Migrating YEAR(2) Columns to YEAR(4)

MySQL 8.0 does not support the `YEAR(2)` data type permitted in older versions of MySQL. Existing `YEAR(2)` columns must be converted to `YEAR(4)` to become usable again. This section provides information about performing that conversion.

Removed YEAR(2) Support in MySQL 8.0

MySQL 8.0 handles `YEAR(2)` columns as follows:

- `YEAR(2)` column definitions for new tables produce an `ER_INVALID_YEAR_COLUMN_LENGTH` error:

```
mysql> CREATE TABLE t1 (y YEAR(2));
ERROR 1818 (HY000): Supports only YEAR or YEAR(4) column.
```

- `YEAR(2)` column in existing tables remain as `YEAR(2)`, but `YEAR(2)` columns in queries produce warnings or errors.
- Several programs or statements convert `YEAR(2)` to `YEAR(4)` automatically:
  - `ALTER TABLE` statements that result in a table rebuild.
  - `REPAIR TABLE` (which `CHECK TABLE` recommends you use if it finds that a table contains `YEAR(2)` columns).
  - `mysql_upgrade` (which uses `REPAIR TABLE`).
  - Dumping with `mysqldump` and reloading the dump file. Unlike the conversions performed by the preceding three items, a dump and reload has the potential to change values.

  A MySQL upgrade usually involves at least one of the last two items. However, with respect to `YEAR(2)`, `mysql_upgrade` is preferable. You should avoid using `mysqldump` because, as noted, that can change values.

Migrating from YEAR(2) to YEAR(4)

To convert `YEAR(2)` columns to `YEAR(4)`, you can do so manually at any time without upgrading. Alternatively, you can upgrade to a version of MySQL with reduced or removed support for `YEAR(2)` (MySQL 5.6.6 or later), then have MySQL convert `YEAR(2)` columns automatically. In the latter case, avoid upgrading by dumping and reloading your data because that can change data values. In addition, if you use replication, there are upgrade considerations you must take into account.

To convert `YEAR(2)` columns to `YEAR(4)` manually, use `ALTER TABLE` or `REPAIR TABLE`. Suppose that a table `t1` has this definition:

```
CREATE TABLE t1 (ycol YEAR(2) NOT NULL DEFAULT '70');
```

Modify the column using `ALTER TABLE` as follows:

```
ALTER TABLE t1 FORCE;
```

The `ALTER TABLE` statement converts the table without changing `YEAR(2)` values. If the server is a replication master, the `ALTER TABLE` statement replicates to slaves and makes the corresponding table change on each one.

Another migration method is to perform a binary upgrade: Install MySQL without dumping and reloading your data. Then run `mysql_upgrade`, which uses `REPAIR TABLE` to convert `YEAR(2)`

columns to YEAR(4) without changing data values. If the server is a replication master, the REPAIR TABLE statements replicate to slaves and make the corresponding table changes on each one, unless you invoke mysql_upgrade with the --skip-write-binlog option.

Upgrades to replication servers usually involve upgrading slaves to a newer version of MySQL, then upgrading the master. For example, if a master and slave both run MySQL 5.5, a typical upgrade sequence involves upgrading the slave to 5.6, then upgrading the master to 5.6. With regard to the different treatment of YEAR(2) as of MySQL 5.6.6, that upgrade sequence results in a problem: Suppose that the slave has been upgraded but not yet the master. Then creating a table containing a YEAR(2) column on the master results in a table containing a YEAR(4) column on the slave. Consequently, these operations will have a different result on the master and slave, if you use statement-based replication:

- Inserting numeric 0. The resulting value has an internal value of 2000 on the master but 0000 on the slave.
- Converting YEAR(2) to string. This operation uses the display value of YEAR(2) on the master but YEAR(4) on the slave.

To avoid such problems, modify all YEAR(2) columns on the master to YEAR(4) before upgrading. (Use ALTER TABLE, as described previously.) Then you can upgrade normally (slave first, then master) without introducing any YEAR(2) to YEAR(4) differences between the master and slave.

One migration method should be avoided: Do not dump your data with mysqldump and reload the dump file after upgrading. This has the potential to change YEAR(2) values, as described previously.

A migration from YEAR(2) to YEAR(4) should also involve examining application code for the possibility of changed behavior under conditions such as these:

- Code that expects selecting a YEAR column to produce exactly two digits.
- Code that does not account for different handling for inserts of numeric 0: Inserting 0 into YEAR(2) or YEAR(4) results in an internal value of 2000 or 0000, respectively.

11.3.5 Automatic Initialization and Updating for TIMESTAMP and DATETIME

TIMESTAMP and DATETIME columns can be automatically initializated and updated to the current date and time (that is, the current timestamp).

For any TIMESTAMP or DATETIME column in a table, you can assign the current timestamp as the default value, the auto-update value, or both:

- An auto-initialized column is set to the current timestamp for inserted rows that specify no value for the column.
- An auto-updated column is automatically updated to the current timestamp when the value of any other column in the row is changed from its current value. An auto-updated column remains unchanged if all other columns are set to their current values. To prevent an auto-updated column from updating when other columns change, explicitly set it to its current value. To update an auto-updated column even when other columns do not change, explicitly set it to the value it should have (for example, set it to CURRENT_TIMESTAMP).

In addition, if the explicit_defaults_for_timestamp system variable is disabled, you can initialize or update any TIMESTAMP (but not DATETIME) column to the current date and time by assigning it a NULL value, unless it has been defined with the NULL attribute to permit NULL values.

To specify automatic properties, use the DEFAULT CURRENT_TIMESTAMP and ON UPDATE CURRENT_TIMESTAMP clauses in column definitions. The order of the clauses does not matter. If both are present in a column definition, either can occur first. Any of the synonyms for CURRENT_TIMESTAMP have the same meaning as CURRENT_TIMESTAMP. These are

`CURRENT_TIMESTAMP()`, `NOW()`, `LOCALTIME`, `LOCALTIME()`, `LOCALTIMESTAMP`, and `LOCALTIMESTAMP()`.

Use of `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP` is specific to `TIMESTAMP` and `DATETIME`. The `DEFAULT` clause also can be used to specify a constant (nonautomatic) default value; for example, `DEFAULT 0` or `DEFAULT '2000-01-01 00:00:00'`.

Note

The following examples use `DEFAULT 0`, a default that can produce warnings or errors depending on whether strict SQL mode or the `NO_ZERO_DATE` SQL mode is enabled. Be aware that the `TRADITIONAL` SQL mode includes strict mode and `NO_ZERO_DATE`. See Section 5.1.8, “Server SQL Modes”.

`TIMESTAMP` or `DATETIME` column definitions can specify the current timestamp for both the default and auto-update values, for one but not the other, or for neither. Different columns can have different combinations of automatic properties. The following rules describe the possibilities:

- With both `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP`, the column has the current timestamp for its default value and is automatically updated to the current timestamp.

```
CREATE TABLE t1 (
  ts TIMESTAMP DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,
  dt DATETIME DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP
);
```

- With a `DEFAULT` clause but no `ON UPDATE CURRENT_TIMESTAMP` clause, the column has the given default value and is not automatically updated to the current timestamp.

  The default depends on whether the `DEFAULT` clause specifies `CURRENT_TIMESTAMP` or a constant value. With `CURRENT_TIMESTAMP`, the default is the current timestamp.

```
CREATE TABLE t1 (
  ts TIMESTAMP DEFAULT CURRENT_TIMESTAMP,
  dt DATETIME DEFAULT CURRENT_TIMESTAMP
);
```

  With a constant, the default is the given value. In this case, the column has no automatic properties at all.

```
CREATE TABLE t1 (
  ts TIMESTAMP DEFAULT 0,
  dt DATETIME DEFAULT 0
);
```

- With an `ON UPDATE CURRENT_TIMESTAMP` clause and a constant `DEFAULT` clause, the column is automatically updated to the current timestamp and has the given constant default value.

```
CREATE TABLE t1 (
  ts TIMESTAMP DEFAULT 0 ON UPDATE CURRENT_TIMESTAMP,
  dt DATETIME DEFAULT 0 ON UPDATE CURRENT_TIMESTAMP
);
```

- With an `ON UPDATE CURRENT_TIMESTAMP` clause but no `DEFAULT` clause, the column is automatically updated to the current timestamp but does not have the current timestamp for its default value.

  The default in this case is type dependent. `TIMESTAMP` has a default of 0 unless defined with the `NULL` attribute, in which case the default is `NULL`.

```
```

```
CREATE TABLE t1 (
  ts1 TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,     -- default 0
  ts2 TIMESTAMP NULL ON UPDATE CURRENT_TIMESTAMP -- default NULL
);
```

`DATETIME` has a default of `NULL` unless defined with the `NOT NULL` attribute, in which case the default is 0.

```
CREATE TABLE t1 (
  dt1 DATETIME ON UPDATE CURRENT_TIMESTAMP,         -- default NULL
  dt2 DATETIME NOT NULL ON UPDATE CURRENT_TIMESTAMP -- default 0
);
```

`TIMESTAMP` and `DATETIME` columns have no automatic properties unless they are specified explicitly, with this exception: If the `explicit_defaults_for_timestamp` system variable is disabled, the *first* `TIMESTAMP` column has both `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP` if neither is specified explicitly. To suppress automatic properties for the first `TIMESTAMP` column, use one of these strategies:

- Enable the `explicit_defaults_for_timestamp` system variable. In this case, the `DEFAULT CURRENT_TIMESTAMP` and `ON UPDATE CURRENT_TIMESTAMP` clauses that specify automatic initialization and updating are available, but are not assigned to any `TIMESTAMP` column unless explicitly included in the column definition.
- Alternatively, if `explicit_defaults_for_timestamp` is disabled, do either of the following:
  - Define the column with a `DEFAULT` clause that specifies a constant default value.
  - Specify the `NULL` attribute. This also causes the column to permit `NULL` values, which means that you cannot assign the current timestamp by setting the column to `NULL`. Assigning `NULL` sets the column to `NULL`, not the current timestamp. To assign the current timestamp, set the column to `CURRENT_TIMESTAMP` or a synonym such as `NOW()`.

Consider these table definitions:

```
CREATE TABLE t1 (
  ts1 TIMESTAMP DEFAULT 0,
  ts2 TIMESTAMP DEFAULT CURRENT_TIMESTAMP
                ON UPDATE CURRENT_TIMESTAMP);
CREATE TABLE t2 (
  ts1 TIMESTAMP NULL,
  ts2 TIMESTAMP DEFAULT CURRENT_TIMESTAMP
                ON UPDATE CURRENT_TIMESTAMP);
CREATE TABLE t3 (
  ts1 TIMESTAMP NULL DEFAULT 0,
  ts2 TIMESTAMP DEFAULT CURRENT_TIMESTAMP
                ON UPDATE CURRENT_TIMESTAMP);
```

The tables have these properties:

- In each table definition, the first `TIMESTAMP` column has no automatic initialization or updating.
- The tables differ in how the `ts1` column handles `NULL` values. For `t1`, `ts1` is `NOT NULL` and assigning it a value of `NULL` sets it to the current timestamp. For `t2` and `t3`, `ts1` permits `NULL` and assigning it a value of `NULL` sets it to `NULL`.
- `t2` and `t3` differ in the default value for `ts1`. For `t2`, `ts1` is defined to permit `NULL`, so the default is also `NULL` in the absence of an explicit `DEFAULT` clause. For `t3`, `ts1` permits `NULL` but has an explicit default of 0.

If a `TIMESTAMP` or `DATETIME` column definition includes an explicit fractional seconds precision value anywhere, the same value must be used throughout the column definition. This is permitted:

```
CREATE TABLE t1 (
  ts TIMESTAMP(6) DEFAULT CURRENT_TIMESTAMP(6) ON UPDATE CURRENT_TIMESTAMP(6)
);
```

This is not permitted:

```
CREATE TABLE t1 (
  ts TIMESTAMP(6) DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP(3)
);
```

TIMESTAMP Initialization and the NULL Attribute

If the `explicit_defaults_for_timestamp` system variable is disabled, `TIMESTAMP` columns by default are `NOT NULL`, cannot contain `NULL` values, and assigning `NULL` assigns the current timestamp. To permit a `TIMESTAMP` column to contain `NULL`, explicitly declare it with the `NULL` attribute. In this case, the default value also becomes `NULL` unless overridden with a `DEFAULT` clause that specifies a different default value. `DEFAULT NULL` can be used to explicitly specify `NULL` as the default value. (For a `TIMESTAMP` column not declared with the `NULL` attribute, `DEFAULT NULL` is invalid.) If a `TIMESTAMP` column permits `NULL` values, assigning `NULL` sets it to `NULL`, not to the current timestamp.

The following table contains several `TIMESTAMP` columns that permit `NULL` values:

```
CREATE TABLE t
(
  ts1 TIMESTAMP NULL DEFAULT NULL,
  ts2 TIMESTAMP NULL DEFAULT 0,
  ts3 TIMESTAMP NULL DEFAULT CURRENT_TIMESTAMP
);
```

A `TIMESTAMP` column that permits `NULL` values does *not* take on the current timestamp at insert time except under one of the following conditions:

- Its default value is defined as `CURRENT_TIMESTAMP` and no value is specified for the column
- `CURRENT_TIMESTAMP` or any of its synonyms such as `NOW()` is explicitly inserted into the column

In other words, a `TIMESTAMP` column defined to permit `NULL` values auto-initializes only if its definition includes `DEFAULT CURRENT_TIMESTAMP`:

```
CREATE TABLE t (ts TIMESTAMP NULL DEFAULT CURRENT_TIMESTAMP);
```

If the `TIMESTAMP` column permits `NULL` values but its definition does not include `DEFAULT CURRENT_TIMESTAMP`, you must explicitly insert a value corresponding to the current date and time. Suppose that tables `t1` and `t2` have these definitions:

```
CREATE TABLE t1 (ts TIMESTAMP NULL DEFAULT '0000-00-00 00:00:00');
CREATE TABLE t2 (ts TIMESTAMP NULL DEFAULT NULL);
```

To set the `TIMESTAMP` column in either table to the current timestamp at insert time, explicitly assign it that value. For example:

```
INSERT INTO t2 VALUES (CURRENT_TIMESTAMP);
INSERT INTO t1 VALUES (NOW());
```

If the `explicit_defaults_for_timestamp` system variable is enabled, `TIMESTAMP` columns permit `NULL` values only if declared with the `NULL` attribute. Also, `TIMESTAMP` columns do not permit assigning `NULL` to assign the current timestamp, whether declared with the `NULL` or `NOT NULL` attribute. To assign the current timestamp, set the column to `CURRENT_TIMESTAMP` or a synonym such as `NOW()`.

11.3.6 Fractional Seconds in Time Values

MySQL 8.0 has fractional seconds support for `TIME`, `DATETIME`, and `TIMESTAMP` values, with up to microseconds (6 digits) precision:

- To define a column that includes a fractional seconds part, use the syntax *`type_name(fsp)`*, where *`type_name`* is `TIME`, `DATETIME`, or `TIMESTAMP`, and *`fsp`* is the fractional seconds precision. For example:

```
CREATE TABLE t1 (t TIME(3), dt DATETIME(6));
```

  The *`fsp`* value, if given, must be in the range 0 to 6. A value of 0 signifies that there is no fractional part. If omitted, the default precision is 0. (This differs from the standard SQL default of 6, for compatibility with previous MySQL versions.)

- Inserting a `TIME`, `DATE`, or `TIMESTAMP` value with a fractional seconds part into a column of the same type but having fewer fractional digits results in rounding, as shown in this example:

```
mysql> CREATE TABLE fractest( c1 TIME(2), c2 DATETIME(2), c3 TIMESTAMP(2) );
Query OK, 0 rows affected (0.33 sec)

mysql> INSERT INTO fractest VALUES
    > ('17:51:04.777', '2014-09-08 17:51:04.777', '2014-09-08 17:51:04.777');
Query OK, 1 row affected (0.03 sec)

mysql> SELECT * FROM fractest;
+-------------+------------------------+------------------------+
| c1          | c2                     | c3                     |
+-------------+------------------------+------------------------+
| 17:51:04.78 | 2014-09-08 17:51:04.78 | 2014-09-08 17:51:04.78 |
+-------------+------------------------+------------------------+
1 row in set (0.00 sec)
```

  No warning or error is given when such rounding occurs. This behavior follows the SQL standard, and is not affected by the server's `sql_mode` setting.

- Functions that take temporal arguments accept values with fractional seconds. Return values from temporal functions include fractional seconds as appropriate. For example, `NOW()` with no argument returns the current date and time with no fractional part, but takes an optional argument from 0 to 6 to specify that the return value includes a fractional seconds part of that many digits.

- Syntax for temporal literals produces temporal values: `DATE 'str'`, `TIME 'str'`, and `TIMESTAMP 'str'`, and the ODBC-syntax equivalents. The resulting value includes a trailing fractional seconds part if specified. Previously, the temporal type keyword was ignored and these constructs produced the string value. See Standard SQL and ODBC Date and Time Literals

11.3.7 Conversion Between Date and Time Types

To some extent, you can convert a value from one temporal type to another. However, there may be some alteration of the value or loss of information. In all cases, conversion between temporal types is subject to the range of valid values for the resulting type. For example, although `DATE`, `DATETIME`, and `TIMESTAMP` values all can be specified using the same set of formats, the types do not all have the same range of values. `TIMESTAMP` values cannot be earlier than `1970` UTC or later than `'2038-01-19 03:14:07'` UTC. This means that a date such as `'1968-01-01'`, while valid as a `DATE` or `DATETIME` value, is not valid as a `TIMESTAMP` value and is converted to `0`.

Conversion of `DATE` values:

- Conversion to a `DATETIME` or `TIMESTAMP` value adds a time part of `'00:00:00'` because the `DATE` value contains no time information.

- Conversion to a `TIME` value is not useful; the result is `'00:00:00'`.

Conversion of DATETIME and TIMESTAMP values:

- Conversion to a DATE value takes fractional seconds into account and rounds the time part. For example, '1999-12-31 23:59:59.499' becomes '1999-12-31', whereas '1999-12-31 23:59:59.500' becomes '2000-01-01'.
- Conversion to a TIME value discards the date part because the TIME type contains no date information.

For conversion of TIME values to other temporal types, the value of CURRENT_DATE() is used for the date part. The TIME is interpreted as elapsed time (not time of day) and added to the date. This means that the date part of the result differs from the current date if the time value is outside the range from '00:00:00' to '23:59:59'.

Suppose that the current date is '2012-01-01'. TIME values of '12:00:00', '24:00:00', and '-12:00:00', when converted to DATETIME or TIMESTAMP values, result in '2012-01-01 12:00:00', '2012-01-02 00:00:00', and '2011-12-31 12:00:00', respectively.

Conversion of TIME to DATE is similar but discards the time part from the result: '2012-01-01', '2012-01-02', and '2011-12-31', respectively.

Explicit conversion can be used to override implicit conversion. For example, in comparison of DATE and DATETIME values, the DATE value is coerced to the DATETIME type by adding a time part of '00:00:00'. To perform the comparison by ignoring the time part of the DATETIME value instead, use the CAST() function in the following way:

```
date_col = CAST(datetime_col AS DATE)
```

Conversion of TIME and DATETIME values to numeric form (for example, by adding +0) depends on whether the value contains a fractional seconds part. TIME(N) or DATETIME(N) is converted to integer when N is 0 (or omitted) and to a DECIMAL value with N decimal digits when N is greater than 0:

```
mysql> SELECT CURTIME(), CURTIME()+0, CURTIME(3)+0;
+-----------+-------------+--------------+
| CURTIME() | CURTIME()+0 | CURTIME(3)+0 |
+-----------+-------------+--------------+
| 09:28:00  |       92800 |    92800.887 |
+-----------+-------------+--------------+
mysql> SELECT NOW(), NOW()+0, NOW(3)+0;
+---------------------+----------------+--------------------+
| NOW()               | NOW()+0        | NOW(3)+0           |
+---------------------+----------------+--------------------+
| 2012-08-15 09:28:00 | 20120815092800 | 20120815092800.889 |
+---------------------+----------------+--------------------+
```

11.3.8 Two-Digit Years in Dates

Date values with two-digit years are ambiguous because the century is unknown. Such values must be interpreted into four-digit form because MySQL stores years internally using four digits.

For DATETIME, DATE, and TIMESTAMP types, MySQL interprets dates specified with ambiguous year values using these rules:

- Year values in the range 00-69 are converted to 2000-2069.
- Year values in the range 70-99 are converted to 1970-1999.

For YEAR, the rules are the same, with this exception: A numeric 00 inserted into YEAR(4) results in 0000 rather than 2000. To specify zero for YEAR(4) and have it be interpreted as 2000, specify it as a string '0' or '00'.

Remember that these rules are only heuristics that provide reasonable guesses as to what your data values mean. If the rules used by MySQL do not produce the values you require, you must provide unambiguous input containing four-digit year values.

`ORDER BY` properly sorts `YEAR` values that have two-digit years.

Some functions like `MIN()` and `MAX()` convert a `YEAR` to a number. This means that a value with a two-digit year does not work properly with these functions. The fix in this case is to convert the `YEAR` to four-digit year format.

11.4 String Types

The string types are `CHAR`, `VARCHAR`, `BINARY`, `VARBINARY`, `BLOB`, `TEXT`, `ENUM`, and `SET`. This section describes how these types work and how to use them in your queries. For string type storage requirements, see Section 11.8, “Data Type Storage Requirements”.

11.4.1 The CHAR and VARCHAR Types

The `CHAR` and `VARCHAR` types are similar, but differ in the way they are stored and retrieved. They also differ in maximum length and in whether trailing spaces are retained.

The `CHAR` and `VARCHAR` types are declared with a length that indicates the maximum number of characters you want to store. For example, `CHAR(30)` can hold up to 30 characters.

The length of a `CHAR` column is fixed to the length that you declare when you create the table. The length can be any value from 0 to 255. When `CHAR` values are stored, they are right-padded with spaces to the specified length. When `CHAR` values are retrieved, trailing spaces are removed unless the `PAD_CHAR_TO_FULL_LENGTH` SQL mode is enabled.

Values in `VARCHAR` columns are variable-length strings. The length can be specified as a value from 0 to 65,535. The effective maximum length of a `VARCHAR` is subject to the maximum row size (65,535 bytes, which is shared among all columns) and the character set used. See Section C.10.4, “Limits on Table Column Count and Row Size”.

In contrast to `CHAR`, `VARCHAR` values are stored as a 1-byte or 2-byte length prefix plus data. The length prefix indicates the number of bytes in the value. A column uses one length byte if values require no more than 255 bytes, two length bytes if values may require more than 255 bytes.

If strict SQL mode is not enabled and you assign a value to a `CHAR` or `VARCHAR` column that exceeds the column's maximum length, the value is truncated to fit and a warning is generated. For truncation of nonspace characters, you can cause an error to occur (rather than a warning) and suppress insertion of the value by using strict SQL mode. See Section 5.1.8, “Server SQL Modes”.

For `VARCHAR` columns, trailing spaces in excess of the column length are truncated prior to insertion and a warning is generated, regardless of the SQL mode in use. For `CHAR` columns, truncation of excess trailing spaces from inserted values is performed silently regardless of the SQL mode.

`VARCHAR` values are not padded when they are stored. Trailing spaces are retained when values are stored and retrieved, in conformance with standard SQL.

The following table illustrates the differences between `CHAR` and `VARCHAR` by showing the result of storing various string values into `CHAR(4)` and `VARCHAR(4)` columns (assuming that the column uses a single-byte character set such as `latin1`).

Value	`CHAR(4)`	Storage Required	`VARCHAR(4)`	Storage Required
`''`	`'    '`	4 bytes	`''`	1 byte
`'ab'`	`'ab  '`	4 bytes	`'ab'`	3 bytes
`'abcd'`	`'abcd'`	4 bytes	`'abcd'`	5 bytes
`'abcdefgh'`	`'abcd'`	4 bytes	`'abcd'`	5 bytes

The values shown as stored in the last row of the table apply *only when not using strict mode*; if MySQL is running in strict mode, values that exceed the column length are *not stored*, and an error results.

`InnoDB` encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored off-page. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.

If a given value is stored into the `CHAR(4)` and `VARCHAR(4)` columns, the values retrieved from the columns are not always the same because trailing spaces are removed from `CHAR` columns upon retrieval. The following example illustrates this difference:

```
mysql> CREATE TABLE vc (v VARCHAR(4), c CHAR(4));
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO vc VALUES ('ab  ', 'ab  ');
Query OK, 1 row affected (0.00 sec)

mysql> SELECT CONCAT('(', v, ')'), CONCAT('(', c, ')') FROM vc;
+---------------------+---------------------+
| CONCAT('(', v, ')') | CONCAT('(', c, ')') |
+---------------------+---------------------+
| (ab  )              | (ab)                |
+---------------------+---------------------+
1 row in set (0.06 sec)
```

Values in `CHAR` and `VARCHAR` columns are sorted and compared according to the character set collation assigned to the column.

Most MySQL collations have a pad attribute of PAD SPACE. The exceptions are Unicode collations based on UCA 9.0.0 and higher, which have a pad attribute of NO PAD. (see Section 10.1.10.1, “Unicode Character Sets”).

To determine the pad attribute for a collation, use the `INFORMATION_SCHEMA COLLATIONS` table, which has a `PAD_ATTRIBUTE` column.

The pad attribute determines how trailing spaces are treated for comparison of nonbinary strings (`CHAR`, `VARCHAR`, and `TEXT` values). NO PAD collations treat spaces at the end of strings like any other character. For PAD SPACE collations, trailing spaces are insignificant in comparisons; strings are compared without regard to any trailing spaces. “Comparison” in this context does not include the `LIKE` pattern-matching operator, for which trailing spaces are significant. For example:

```
mysql> CREATE TABLE names (myname CHAR(10));
Query OK, 0 rows affected (0.03 sec)

mysql> INSERT INTO names VALUES ('Monty');
Query OK, 1 row affected (0.00 sec)

mysql> SELECT myname = 'Monty', myname = 'Monty  ' FROM names;
+------------------+--------------------+
| myname = 'Monty' | myname = 'Monty  ' |
+------------------+--------------------+
|                1 |                  1 |
+------------------+--------------------+
1 row in set (0.00 sec)

mysql> SELECT myname LIKE 'Monty', myname LIKE 'Monty  ' FROM names;
+--------------------+----------------------+
| myname LIKE 'Monty' | myname LIKE 'Monty  ' |
+--------------------+----------------------+
|                  1 |                    0 |
+--------------------+----------------------+
1 row in set (0.00 sec)
```

This is true for all MySQL versions, and is not affected by the server SQL mode.

Note

For more information about MySQL character sets and collations, see Section 10.1, “Character Set Support”. For additional information about storage requirements, see Section 11.8, “Data Type Storage Requirements”.

For those cases where trailing pad characters are stripped or comparisons ignore them, if a column has an index that requires unique values, inserting into the column values that differ only in number of trailing pad characters will result in a duplicate-key error. For example, if a table contains `'a'`, an attempt to store `'a '` causes a duplicate-key error.

11.4.2 The BINARY and VARBINARY Types

The `BINARY` and `VARBINARY` types are similar to `CHAR` and `VARCHAR`, except that they contain binary strings rather than nonbinary strings. That is, they contain byte strings rather than character strings. This means they have the `binary` character set and collation, and comparison and sorting are based on the numeric values of the bytes in the values.

The permissible maximum length is the same for `BINARY` and `VARBINARY` as it is for `CHAR` and `VARCHAR`, except that the length for `BINARY` and `VARBINARY` is a length in bytes rather than in characters.

The `BINARY` and `VARBINARY` data types are distinct from the `CHAR BINARY` and `VARCHAR BINARY` data types. For the latter types, the `BINARY` attribute does not cause the column to be treated as a binary string column. Instead, it causes the binary (`_bin`) collation for the column character set to be used, and the column itself contains nonbinary character strings rather than binary byte strings. For example, `CHAR(5) BINARY` is treated as `CHAR(5) CHARACTER SET utf8mb4 COLLATE utf8mb4_bin`, assuming that the default character set is `utf8mb4`. This differs from `BINARY(5)`, which stores 5-bytes binary strings that have the `binary` character set and collation. For information about differences between binary strings and binary collations for nonbinary strings, see Section 10.1.8.5, “The binary Collation Compared to _bin Collations”.

If strict SQL mode is not enabled and you assign a value to a `BINARY` or `VARBINARY` column that exceeds the column's maximum length, the value is truncated to fit and a warning is generated. For cases of truncation, you can cause an error to occur (rather than a warning) and suppress insertion of the value by using strict SQL mode. See Section 5.1.8, “Server SQL Modes”.

When `BINARY` values are stored, they are right-padded with the pad value to the specified length. The pad value is `0x00` (the zero byte). Values are right-padded with `0x00` on insert, and no trailing bytes are removed on select. All bytes are significant in comparisons, including `ORDER BY` and `DISTINCT` operations. `0x00` bytes and spaces are different in comparisons, with `0x00` < space.

Example: For a `BINARY(3)` column, `'a '` becomes `'a \0'` when inserted. `'a\0'` becomes `'a \0\0'` when inserted. Both inserted values remain unchanged when selected.

For `VARBINARY`, there is no padding on insert and no bytes are stripped on select. All bytes are significant in comparisons, including `ORDER BY` and `DISTINCT` operations. `0x00` bytes and spaces are different in comparisons, with `0x00` < space.

For those cases where trailing pad bytes are stripped or comparisons ignore them, if a column has an index that requires unique values, inserting into the column values that differ only in number of trailing pad bytes will result in a duplicate-key error. For example, if a table contains `'a'`, an attempt to store `'a\0'` causes a duplicate-key error.

You should consider the preceding padding and stripping characteristics carefully if you plan to use the `BINARY` data type for storing binary data and you require that the value retrieved be exactly the same as the value stored. The following example illustrates how `0x00`-padding of `BINARY` values affects column value comparisons:

```
mysql> CREATE TABLE t (c BINARY(3));
```

```
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO t SET c = 'a';
Query OK, 1 row affected (0.01 sec)

mysql> SELECT HEX(c), c = 'a', c = 'a\0\0' from t;
+--------+---------+-------------+
| HEX(c) | c = 'a' | c = 'a\0\0' |
+--------+---------+-------------+
| 610000 |       0 |           1 |
+--------+---------+-------------+
1 row in set (0.09 sec)
```

If the value retrieved must be the same as the value specified for storage with no padding, it might be preferable to use `VARBINARY` or one of the `BLOB` data types instead.

11.4.3 The BLOB and TEXT Types

A `BLOB` is a binary large object that can hold a variable amount of data. The four `BLOB` types are `TINYBLOB`, `BLOB`, `MEDIUMBLOB`, and `LONGBLOB`. These differ only in the maximum length of the values they can hold. The four `TEXT` types are `TINYTEXT`, `TEXT`, `MEDIUMTEXT`, and `LONGTEXT`. These correspond to the four `BLOB` types and have the same maximum lengths and storage requirements. See Section 11.8, “Data Type Storage Requirements”.

`BLOB` values are treated as binary strings (byte strings). They have the `binary` character set and collation, and comparison and sorting are based on the numeric values of the bytes in column values. `TEXT` values are treated as nonbinary strings (character strings). They have a character set other than `binary`, and values are sorted and compared based on the collation of the character set.

If strict SQL mode is not enabled and you assign a value to a `BLOB` or `TEXT` column that exceeds the column's maximum length, the value is truncated to fit and a warning is generated. For truncation of nonspace characters, you can cause an error to occur (rather than a warning) and suppress insertion of the value by using strict SQL mode. See Section 5.1.8, “Server SQL Modes”.

Truncation of excess trailing spaces from values to be inserted into `TEXT` columns always generates a warning, regardless of the SQL mode.

For `TEXT` and `BLOB` columns, there is no padding on insert and no bytes are stripped on select.

If a `TEXT` column is indexed, index entry comparisons are space-padded at the end. This means that, if the index requires unique values, duplicate-key errors will occur for values that differ only in the number of trailing spaces. For example, if a table contains `'a'`, an attempt to store `'a '` causes a duplicate-key error. This is not true for `BLOB` columns.

In most respects, you can regard a `BLOB` column as a `VARBINARY` column that can be as large as you like. Similarly, you can regard a `TEXT` column as a `VARCHAR` column. `BLOB` and `TEXT` differ from `VARBINARY` and `VARCHAR` in the following ways:

- For indexes on `BLOB` and `TEXT` columns, you must specify an index prefix length. For `CHAR` and `VARCHAR`, a prefix length is optional. See Section 8.3.4, “Column Indexes”.
- `BLOB` and `TEXT` columns cannot have `DEFAULT` values.

If you use the `BINARY` attribute with a `TEXT` data type, the column is assigned the binary (`_bin`) collation of the column character set.

`LONG` and `LONG VARCHAR` map to the `MEDIUMTEXT` data type. This is a compatibility feature.

MySQL Connector/ODBC defines `BLOB` values as `LONGVARBINARY` and `TEXT` values as `LONGVARCHAR`.

Because `BLOB` and `TEXT` values can be extremely long, you might encounter some constraints in using them:

- Only the first `max_sort_length` bytes of the column are used when sorting. The default value of `max_sort_length` is 1024. You can make more bytes significant in sorting or grouping by increasing the value of `max_sort_length` at server startup or runtime. Any client can change the value of its session `max_sort_length` variable:

```
mysql> SET max_sort_length = 2000;
mysql> SELECT id, comment FROM t
    -> ORDER BY comment;
```

- Instances of `BLOB` or `TEXT` columns in the result of a query that is processed using a temporary table causes the server to use a table on disk rather than in memory because the `MEMORY` storage engine does not support those data types (see Section 8.4.4, “Internal Temporary Table Use in MySQL”). Use of disk incurs a performance penalty, so include `BLOB` or `TEXT` columns in the query result only if they are really needed. For example, avoid using `SELECT *`, which selects all columns.

- The maximum size of a `BLOB` or `TEXT` object is determined by its type, but the largest value you actually can transmit between the client and server is determined by the amount of available memory and the size of the communications buffers. You can change the message buffer size by changing the value of the `max_allowed_packet` variable, but you must do so for both the server and your client program. For example, both `mysql` and `mysqldump` enable you to change the client-side `max_allowed_packet` value. See Section 5.1.1, “Configuring the Server”, Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”, and Section 4.5.4, “`mysqldump` — A Database Backup Program”. You may also want to compare the packet sizes and the size of the data objects you are storing with the storage requirements, see Section 11.8, “Data Type Storage Requirements”

Each `BLOB` or `TEXT` value is represented internally by a separately allocated object. This is in contrast to all other data types, for which storage is allocated once per column when the table is opened.

In some cases, it may be desirable to store binary data such as media files in `BLOB` or `TEXT` columns. You may find MySQL's string handling functions useful for working with such data. See Section 12.5, “String Functions”. For security and other reasons, it is usually preferable to do so using application code rather than giving application users the `FILE` privilege. You can discuss specifics for various languages and platforms in the MySQL Forums (http://forums.mysql.com/).

11.4.4 The ENUM Type

An `ENUM` is a string object with a value chosen from a list of permitted values that are enumerated explicitly in the column specification at table creation time. It has these advantages:

- Compact data storage in situations where a column has a limited set of possible values. The strings you specify as input values are automatically encoded as numbers. See Section 11.8, “Data Type Storage Requirements” for the storage requirements for `ENUM` types.

- Readable queries and output. The numbers are translated back to the corresponding strings in query results.

and these potential issues to consider:

- If you make enumeration values that look like numbers, it is easy to mix up the literal values with their internal index numbers, as explained in Enumeration Limitations.

- Using `ENUM` columns in `ORDER BY` clauses requires extra care, as explained in Enumeration Sorting.

- Creating and Using ENUM Columns

- Index Values for Enumeration Literals

- Handling of Enumeration Literals

- Empty or NULL Enumeration Values

- Enumeration Sorting

- Enumeration Limitations

Creating and Using ENUM Columns

An enumeration value must be a quoted string literal. For example, you can create a table with an `ENUM` column like this:

```
CREATE TABLE shirts (
    name VARCHAR(40),
    size ENUM('x-small', 'small', 'medium', 'large', 'x-large')
);
INSERT INTO shirts (name, size) VALUES ('dress shirt','large'), ('t-shirt','medium'),
  ('polo shirt','small');
SELECT name, size FROM shirts WHERE size = 'medium';
+---------+--------+
| name    | size   |
+---------+--------+
| t-shirt | medium |
+---------+--------+
UPDATE shirts SET size = 'small' WHERE size = 'large';
COMMIT;
```

Inserting 1 million rows into this table with a value of `'medium'` would require 1 million bytes of storage, as opposed to 6 million bytes if you stored the actual string `'medium'` in a `VARCHAR` column.

Index Values for Enumeration Literals

Each enumeration value has an index:

- The elements listed in the column specification are assigned index numbers, beginning with 1.
- The index value of the empty string error value is 0. This means that you can use the following `SELECT` statement to find rows into which invalid `ENUM` values were assigned:

  ```
  mysql> SELECT * FROM tbl_name WHERE enum_col=0;
  ```

- The index of the `NULL` value is `NULL`.
- The term “index” here refers to a position within the list of enumeration values. It has nothing to do with table indexes.

For example, a column specified as `ENUM('Mercury', 'Venus', 'Earth')` can have any of the values shown here. The index of each value is also shown.

Value	Index
`NULL`	`NULL`
`''`	0
`'Mercury'`	1
`'Venus'`	2
`'Earth'`	3

An `ENUM` column can have a maximum of 65,535 distinct elements.

If you retrieve an `ENUM` value in a numeric context, the column value's index is returned. For example, you can retrieve numeric values from an `ENUM` column like this:

```
mysql> SELECT enum_col+0 FROM tbl_name;
```

Functions such as `SUM()` or `AVG()` that expect a numeric argument cast the argument to a number if necessary. For `ENUM` values, the index number is used in the calculation.

Handling of Enumeration Literals

Trailing spaces are automatically deleted from `ENUM` member values in the table definition when a table is created.

When retrieved, values stored into an `ENUM` column are displayed using the lettercase that was used in the column definition. Note that `ENUM` columns can be assigned a character set and collation. For binary or case-sensitive collations, lettercase is taken into account when assigning values to the column.

If you store a number into an `ENUM` column, the number is treated as the index into the possible values, and the value stored is the enumeration member with that index. (However, this does *not* work with `LOAD DATA`, which treats all input as strings.) If the numeric value is quoted, it is still interpreted as an index if there is no matching string in the list of enumeration values. For these reasons, it is not advisable to define an `ENUM` column with enumeration values that look like numbers, because this can easily become confusing. For example, the following column has enumeration members with string values of `'0'`, `'1'`, and `'2'`, but numeric index values of `1`, `2`, and `3`:

```
numbers ENUM('0','1','2')
```

If you store `2`, it is interpreted as an index value, and becomes `'1'` (the value with index 2). If you store `'2'`, it matches an enumeration value, so it is stored as `'2'`. If you store `'3'`, it does not match any enumeration value, so it is treated as an index and becomes `'2'` (the value with index 3).

```
mysql> INSERT INTO t (numbers) VALUES(2),('2'),('3');
mysql> SELECT * FROM t;
+---------+
| numbers |
+---------+
| 1       |
| 2       |
| 2       |
+---------+
```

To determine all possible values for an `ENUM` column, use `SHOW COLUMNS FROM tbl_name LIKE 'enum_col'` and parse the `ENUM` definition in the `Type` column of the output.

In the C API, `ENUM` values are returned as strings. For information about using result set metadata to distinguish them from other strings, see Section 27.7.5, “C API Data Structures”.

Empty or NULL Enumeration Values

An enumeration value can also be the empty string (`''`) or `NULL` under certain circumstances:

- If you insert an invalid value into an `ENUM` (that is, a string not present in the list of permitted values), the empty string is inserted instead as a special error value. This string can be distinguished from a “normal” empty string by the fact that this string has the numeric value 0. See Index Values for Enumeration Literals for details about the numeric indexes for the enumeration values.

  If strict SQL mode is enabled, attempts to insert invalid `ENUM` values result in an error.

- If an `ENUM` column is declared to permit `NULL`, the `NULL` value is a valid value for the column, and the default value is `NULL`. If an `ENUM` column is declared `NOT NULL`, its default value is the first element of the list of permitted values.

Enumeration Sorting

`ENUM` values are sorted based on their index numbers, which depend on the order in which the enumeration members were listed in the column specification. For example, `'b'` sorts before `'a'` for `ENUM('b', 'a')`. The empty string sorts before nonempty strings, and `NULL` values sort before all other enumeration values.

To prevent unexpected results when using the `ORDER BY` clause on an `ENUM` column, use one of these techniques:

- Specify the `ENUM` list in alphabetic order.
- Make sure that the column is sorted lexically rather than by index number by coding `ORDER BY CAST(col AS CHAR)` or `ORDER BY CONCAT(col)`.

Enumeration Limitations

An enumeration value cannot be an expression, even one that evaluates to a string value.

For example, this `CREATE TABLE` statement does *not* work because the `CONCAT` function cannot be used to construct an enumeration value:

```
CREATE TABLE sizes (
    size ENUM('small', CONCAT('med','ium'), 'large')
);
```

You also cannot employ a user variable as an enumeration value. This pair of statements do *not* work:

```
SET @mysize = 'medium';

CREATE TABLE sizes (
    size ENUM('small', @mysize, 'large')
);
```

We strongly recommend that you do *not* use numbers as enumeration values, because it does not save on storage over the appropriate `TINYINT` or `SMALLINT` type, and it is easy to mix up the strings and the underlying number values (which might not be the same) if you quote the `ENUM` values incorrectly. If you do use a number as an enumeration value, always enclose it in quotation marks. If the quotation marks are omitted, the number is regarded as an index. See Handling of Enumeration Literals to see how even a quoted number could be mistakenly used as a numeric index value.

Duplicate values in the definition cause a warning, or an error if strict SQL mode is enabled.

11.4.5 The SET Type

A `SET` is a string object that can have zero or more values, each of which must be chosen from a list of permitted values specified when the table is created. `SET` column values that consist of multiple set members are specified with members separated by commas (`,`). A consequence of this is that `SET` member values should not themselves contain commas.

For example, a column specified as `SET('one', 'two') NOT NULL` can have any of these values:

```
''
'one'
'two'
'one,two'
```

A `SET` column can have a maximum of 64 distinct members.

Duplicate values in the definition cause a warning, or an error if strict SQL mode is enabled.

Trailing spaces are automatically deleted from `SET` member values in the table definition when a table is created.

When retrieved, values stored in a `SET` column are displayed using the lettercase that was used in the column definition. Note that `SET` columns can be assigned a character set and collation. For binary or case-sensitive collations, lettercase is taken into account when assigning values to the column.

MySQL stores `SET` values numerically, with the low-order bit of the stored value corresponding to the first set member. If you retrieve a `SET` value in a numeric context, the value retrieved has bits set corresponding to the set members that make up the column value. For example, you can retrieve numeric values from a `SET` column like this:

```
mysql> SELECT set_col+0 FROM tbl_name;
```

If a number is stored into a `SET` column, the bits that are set in the binary representation of the number determine the set members in the column value. For a column specified as `SET('a','b','c','d')`, the members have the following decimal and binary values.

SET Member	Decimal Value	Binary Value
`'a'`	`1`	`0001`
`'b'`	`2`	`0010`
`'c'`	`4`	`0100`
`'d'`	`8`	`1000`

If you assign a value of `9` to this column, that is `1001` in binary, so the first and fourth `SET` value members `'a'` and `'d'` are selected and the resulting value is `'a,d'`.

For a value containing more than one `SET` element, it does not matter what order the elements are listed in when you insert the value. It also does not matter how many times a given element is listed in the value. When the value is retrieved later, each element in the value appears once, with elements listed according to the order in which they were specified at table creation time. For example, suppose that a column is specified as `SET('a','b','c','d')`:

```
mysql> CREATE TABLE myset (col SET('a', 'b', 'c', 'd'));
```

If you insert the values `'a,d'`, `'d,a'`, `'a,d,d'`, `'a,d,a'`, and `'d,a,d'`:

```
mysql> INSERT INTO myset (col) VALUES
-> ('a,d'), ('d,a'), ('a,d,a'), ('a,d,d'), ('d,a,d');
Query OK, 5 rows affected (0.01 sec)
Records: 5  Duplicates: 0  Warnings: 0
```

Then all these values appear as `'a,d'` when retrieved:

```
mysql> SELECT col FROM myset;
+------+
| col  |
+------+
| a,d  |
| a,d  |
| a,d  |
| a,d  |
| a,d  |
+------+
5 rows in set (0.04 sec)
```

If you set a `SET` column to an unsupported value, the value is ignored and a warning is issued:

```
mysql> INSERT INTO myset (col) VALUES ('a,d,d,s');
Query OK, 1 row affected, 1 warning (0.03 sec)

mysql> SHOW WARNINGS;
+---------+------+-------------------------------------------+
| Level   | Code | Message                                   |
+---------+------+-------------------------------------------+
```

```
| Warning | 1265 | Data truncated for column 'col' at row 1 |
+---------+------+------------------------------------------+
1 row in set (0.04 sec)

mysql> SELECT col FROM myset;
+------+
| col  |
+------+
| a,d  |
| a,d  |
| a,d  |
| a,d  |
| a,d  |
| a,d  |
+------+
6 rows in set (0.01 sec)
```

If strict SQL mode is enabled, attempts to insert invalid `SET` values result in an error.

`SET` values are sorted numerically. `NULL` values sort before non-`NULL` `SET` values.

Functions such as `SUM()` or `AVG()` that expect a numeric argument cast the argument to a number if necessary. For `SET` values, the cast operation causes the numeric value to be used.

Normally, you search for `SET` values using the `FIND_IN_SET()` function or the `LIKE` operator:

```
mysql> SELECT * FROM tbl_name WHERE FIND_IN_SET('value',set_col)>0;
mysql> SELECT * FROM tbl_name WHERE set_col LIKE '%value%';
```

The first statement finds rows where *`set_col`* contains the *`value`* set member. The second is similar, but not the same: It finds rows where *`set_col`* contains *`value`* anywhere, even as a substring of another set member.

The following statements also are permitted:

```
mysql> SELECT * FROM tbl_name WHERE set_col & 1;
mysql> SELECT * FROM tbl_name WHERE set_col = 'val1,val2';
```

The first of these statements looks for values containing the first set member. The second looks for an exact match. Be careful with comparisons of the second type. Comparing set values to *`'val1,val2'`* returns different results than comparing values to *`'val2,val1'`*. You should specify the values in the same order they are listed in the column definition.

To determine all possible values for a `SET` column, use `SHOW COLUMNS FROM` *`tbl_name`* `LIKE` *`set_col`* and parse the `SET` definition in the `Type` column of the output.

In the C API, `SET` values are returned as strings. For information about using result set metadata to distinguish them from other strings, see Section 27.7.5, “C API Data Structures”.

11.5 Extensions for Spatial Data

The Open Geospatial Consortium (OGC) is an international consortium of more than 250 companies, agencies, and universities participating in the development of publicly available conceptual solutions that can be useful with all kinds of applications that manage spatial data.

The Open Geospatial Consortium publishes the *OpenGIS® Implementation Standard for Geographic information - Simple feature access - Part 2: SQL option*, a document that proposes several conceptual ways for extending an SQL RDBMS to support spatial data. This specification is available from the OGC Web site at http://www.opengeospatial.org/standards/sfs.

Following the OGC specification, MySQL implements spatial extensions as a subset of the **SQL with Geometry Types** environment. This term refers to an SQL environment that has been extended

with a set of geometry types. A geometry-valued SQL column is implemented as a column that has a geometry type. The specification describes a set of SQL geometry types, as well as functions on those types to create and analyze geometry values.

MySQL spatial extensions enable the generation, storage, and analysis of geographic features:

- Data types for representing spatial values
- Functions for manipulating spatial values
- Spatial indexing for improved access times to spatial columns

The spatial data types and functions are available for `MyISAM`, `InnoDB`, and `ARCHIVE` tables. For indexing spatial columns, `MyISAM` and `InnoDB` support both `SPATIAL` and non-`SPATIAL` indexes. The other storage engines support non-`SPATIAL` indexes, as described in Section 13.1.12, “CREATE INDEX Syntax”.

A **geographic feature** is anything in the world that has a location. A feature can be:

- An entity. For example, a mountain, a pond, a city.
- A space. For example, town district, the tropics.
- A definable location. For example, a crossroad, as a particular place where two streets intersect.

Some documents use the term **geospatial feature** to refer to geographic features.

Geometry is another word that denotes a geographic feature. Originally the word **geometry** meant measurement of the earth. Another meaning comes from cartography, referring to the geometric features that cartographers use to map the world.

The discussion here considers these terms synonymous: **geographic feature**, **geospatial feature**, **feature**, or **geometry**. The term most commonly used is **geometry**, defined as *a point or an aggregate of points representing anything in the world that has a location*.

The following material covers these topics:

- The spatial data types implemented in MySQL model
- The basis of the spatial extensions in the OpenGIS geometry model
- Data formats for representing spatial data
- How to use spatial data in MySQL
- Use of indexing for spatial data
- MySQL differences from the OpenGIS specification

For information about functions that operate on spatial data, see Section 12.15, “Spatial Analysis Functions”.

MySQL GIS Conformance and Compatibility

MySQL does not implement the following GIS features:

- Additional Metadata Views

  OpenGIS specifications propose several additional metadata views. For example, a system view named `GEOMETRY_COLUMNS` contains a description of geometry columns, one row for each geometry column in the database.

- The OpenGIS function `Length()` on `LineString` and `MultiLineString` should be called in MySQL as `ST_Length()`

  The problem is that there is an existing SQL function `Length()` that calculates the length of string values, and sometimes it is not possible to distinguish whether the function is called in a textual or spatial context.

Additional Resources

- The Open Geospatial Consortium publishes the *OpenGIS® Implementation Standard for Geographic information - Simple feature access - Part 2: SQL option*, a document that proposes several conceptual ways for extending an SQL RDBMS to support spatial data. The Open Geospatial Consortium (OGC) maintains a Web site at http://www.opengeospatial.org/. The specification is available there at http://www.opengeospatial.org/standards/sfs. It contains additional information relevant to the material here.
- If you have questions or concerns about the use of the spatial extensions to MySQL, you can discuss them in the GIS forum: http://forums.mysql.com/list.php?23.

11.5.1 Spatial Data Types

MySQL has data types that correspond to OpenGIS classes. Some of these types hold single geometry values:

- `GEOMETRY`
- `POINT`
- `LINESTRING`
- `POLYGON`

`GEOMETRY` can store geometry values of any type. The other single-value types (`POINT`, `LINESTRING`, and `POLYGON`) restrict their values to a particular geometry type.

The other data types hold collections of values:

- `MULTIPOINT`
- `MULTILINESTRING`
- `MULTIPOLYGON`
- `GEOMETRYCOLLECTION`

`GEOMETRYCOLLECTION` can store a collection of objects of any type. The other collection types (`MULTIPOINT`, `MULTILINESTRING`, `MULTIPOLYGON`, and `GEOMETRYCOLLECTION`) restrict collection members to those having a particular geometry type.

MySQL spatial data types have their basis in the OpenGIS geometry model, described in Section 11.5.2, “The OpenGIS Geometry Model”. For examples showing how to use spatial data types in MySQL, see Section 11.5.5, “Creating Spatial Columns”.

11.5.2 The OpenGIS Geometry Model

The set of geometry types proposed by OGC's **SQL with Geometry Types** environment is based on the **OpenGIS Geometry Model**. In this model, each geometric object has the following general properties:

- It is associated with a spatial reference system, which describes the coordinate space in which the object is defined.

- It belongs to some geometry class.

11.5.2.1 The Geometry Class Hierarchy

The geometry classes define a hierarchy as follows:

- `Geometry` (noninstantiable)
  - `Point` (instantiable)
  - `Curve` (noninstantiable)
    - `LineString` (instantiable)
      - `Line`
      - `LinearRing`
  - `Surface` (noninstantiable)
    - `Polygon` (instantiable)
  - `GeometryCollection` (instantiable)
    - `MultiPoint` (instantiable)
    - `MultiCurve` (noninstantiable)
      - `MultiLineString` (instantiable)
    - `MultiSurface` (noninstantiable)
      - `MultiPolygon` (instantiable)

It is not possible to create objects in noninstantiable classes. It is possible to create objects in instantiable classes. All classes have properties, and instantiable classes may also have assertions (rules that define valid class instances).

`Geometry` is the base class. It is an abstract class. The instantiable subclasses of `Geometry` are restricted to zero-, one-, and two-dimensional geometric objects that exist in two-dimensional coordinate space. All instantiable geometry classes are defined so that valid instances of a geometry class are topologically closed (that is, all defined geometries include their boundary).

The base `Geometry` class has subclasses for `Point`, `Curve`, `Surface`, and `GeometryCollection`:

- `Point` represents zero-dimensional objects.
- `Curve` represents one-dimensional objects, and has subclass `LineString`, with sub-subclasses `Line` and `LinearRing`.
- `Surface` is designed for two-dimensional objects and has subclass `Polygon`.
- `GeometryCollection` has specialized zero-, one-, and two-dimensional collection classes named `MultiPoint`, `MultiLineString`, and `MultiPolygon` for modeling geometries corresponding to collections of `Points`, `LineStrings`, and `Polygons`, respectively. `MultiCurve` and `MultiSurface` are introduced as abstract superclasses that generalize the collection interfaces to handle `Curves` and `Surfaces`.

`Geometry`, `Curve`, `Surface`, `MultiCurve`, and `MultiSurface` are defined as noninstantiable classes. They define a common set of methods for their subclasses and are included for extensibility.

`Point`, `LineString`, `Polygon`, `GeometryCollection`, `MultiPoint`, `MultiLineString`, and `MultiPolygon` are instantiable classes.

11.5.2.2 Geometry Class

`Geometry` is the root class of the hierarchy. It is a noninstantiable class but has a number of properties, described in the following list, that are common to all geometry values created from any of the `Geometry` subclasses. Particular subclasses have their own specific properties, described later.

Geometry Properties

A geometry value has the following properties:

- Its **type**. Each geometry belongs to one of the instantiable classes in the hierarchy.
- Its **SRID**, or spatial reference identifier. This value identifies the geometry's associated spatial reference system that describes the coordinate space in which the geometry object is defined.

  In MySQL, the SRID value is an integer associated with the geometry value. The maximum usable SRID value is $2^{32}-1$. If a larger value is given, only the lower 32 bits are used. All computations are done assuming SRID 0, regardless of the actual SRID value. SRID 0 represents an infinite flat Cartesian plane with no units assigned to its axes.
- Its **coordinates** in its spatial reference system, represented as double-precision (8-byte) numbers. All nonempty geometries include at least one pair of (X,Y) coordinates. Empty geometries contain no coordinates.

  Coordinates are related to the SRID. For example, in different coordinate systems, the distance between two objects may differ even when objects have the same coordinates, because the distance on the **planar** coordinate system and the distance on the **geodetic** system (coordinates on the Earth's surface) are different things.
- Its **interior**, **boundary**, and **exterior**.

  Every geometry occupies some position in space. The exterior of a geometry is all space not occupied by the geometry. The interior is the space occupied by the geometry. The boundary is the interface between the geometry's interior and exterior.
- Its **MBR** (minimum bounding rectangle), or envelope. This is the bounding geometry, formed by the minimum and maximum (X,Y) coordinates:

  ```
  ((MINX MINY, MAXX MINY, MAXX MAXY, MINX MAXY, MINX MINY))
  ```
- Whether the value is **simple** or **nonsimple**. Geometry values of types (`LineString`, `MultiPoint`, `MultiLineString`) are either simple or nonsimple. Each type determines its own assertions for being simple or nonsimple.
- Whether the value is **closed** or **not closed**. Geometry values of types (`LineString`, `MultiString`) are either closed or not closed. Each type determines its own assertions for being closed or not closed.
- Whether the value is **empty** or **nonempty** A geometry is empty if it does not have any points. Exterior, interior, and boundary of an empty geometry are not defined (that is, they are represented by a `NULL` value). An empty geometry is defined to be always simple and has an area of 0.
- Its **dimension**. A geometry can have a dimension of −1, 0, 1, or 2:
  - −1 for an empty geometry.
  - 0 for a geometry with no length and no area.
  - 1 for a geometry with nonzero length and zero area.
  - 2 for a geometry with nonzero area.

`Point` objects have a dimension of zero. `LineString` objects have a dimension of 1. `Polygon` objects have a dimension of 2. The dimensions of `MultiPoint`, `MultiLineString`, and `MultiPolygon` objects are the same as the dimensions of the elements they consist of.

11.5.2.3 Point Class

A `Point` is a geometry that represents a single location in coordinate space.

`Point` Examples

- Imagine a large-scale map of the world with many cities. A `Point` object could represent each city.
- On a city map, a `Point` object could represent a bus stop.

`Point` Properties

- X-coordinate value.
- Y-coordinate value.
- `Point` is defined as a zero-dimensional geometry.
- The boundary of a `Point` is the empty set.

11.5.2.4 Curve Class

A `Curve` is a one-dimensional geometry, usually represented by a sequence of points. Particular subclasses of `Curve` define the type of interpolation between points. `Curve` is a noninstantiable class.

`Curve` Properties

- A `Curve` has the coordinates of its points.
- A `Curve` is defined as a one-dimensional geometry.
- A `Curve` is simple if it does not pass through the same point twice, with the exception that a curve can still be simple if the start and end points are the same.
- A `Curve` is closed if its start point is equal to its endpoint.
- The boundary of a closed `Curve` is empty.
- The boundary of a nonclosed `Curve` consists of its two endpoints.
- A `Curve` that is simple and closed is a `LinearRing`.

11.5.2.5 LineString Class

A `LineString` is a `Curve` with linear interpolation between points.

`LineString` Examples

- On a world map, `LineString` objects could represent rivers.
- In a city map, `LineString` objects could represent streets.

`LineString` Properties

- A `LineString` has coordinates of segments, defined by each consecutive pair of points.
- A `LineString` is a `Line` if it consists of exactly two points.
- A `LineString` is a `LinearRing` if it is both closed and simple.

11.5.2.6 Surface Class

A `Surface` is a two-dimensional geometry. It is a noninstantiable class. Its only instantiable subclass is `Polygon`.

`Surface` Properties

- A `Surface` is defined as a two-dimensional geometry.
- The OpenGIS specification defines a simple `Surface` as a geometry that consists of a single “patch” that is associated with a single exterior boundary and zero or more interior boundaries.
- The boundary of a simple `Surface` is the set of closed curves corresponding to its exterior and interior boundaries.

11.5.2.7 Polygon Class

A `Polygon` is a planar `Surface` representing a multisided geometry. It is defined by a single exterior boundary and zero or more interior boundaries, where each interior boundary defines a hole in the `Polygon`.

`Polygon` Examples

- On a region map, `Polygon` objects could represent forests, districts, and so on.

`Polygon` Assertions

- The boundary of a `Polygon` consists of a set of `LinearRing` objects (that is, `LineString` objects that are both simple and closed) that make up its exterior and interior boundaries.
- A `Polygon` has no rings that cross. The rings in the boundary of a `Polygon` may intersect at a `Point`, but only as a tangent.
- A `Polygon` has no lines, spikes, or punctures.
- A `Polygon` has an interior that is a connected point set.
- A `Polygon` may have holes. The exterior of a `Polygon` with holes is not connected. Each hole defines a connected component of the exterior.

The preceding assertions make a `Polygon` a simple geometry.

11.5.2.8 GeometryCollection Class

A `GeometryCollection` is a geometry that is a collection of one or more geometries of any class.

All the elements in a `GeometryCollection` must be in the same spatial reference system (that is, in the same coordinate system). There are no other constraints on the elements of a `GeometryCollection`, although the subclasses of `GeometryCollection` described in the following sections may restrict membership. Restrictions may be based on:

- Element type (for example, a `MultiPoint` may contain only `Point` elements)
- Dimension
- Constraints on the degree of spatial overlap between elements

11.5.2.9 MultiPoint Class

A `MultiPoint` is a geometry collection composed of `Point` elements. The points are not connected or ordered in any way.

`MultiPoint` Examples

- On a world map, a `MultiPoint` could represent a chain of small islands.
- On a city map, a `MultiPoint` could represent the outlets for a ticket office.

`MultiPoint` Properties

- A `MultiPoint` is a zero-dimensional geometry.
- A `MultiPoint` is simple if no two of its `Point` values are equal (have identical coordinate values).
- The boundary of a `MultiPoint` is the empty set.

11.5.2.10 MultiCurve Class

A `MultiCurve` is a geometry collection composed of `Curve` elements. `MultiCurve` is a noninstantiable class.

`MultiCurve` Properties

- A `MultiCurve` is a one-dimensional geometry.
- A `MultiCurve` is simple if and only if all of its elements are simple; the only intersections between any two elements occur at points that are on the boundaries of both elements.
- A `MultiCurve` boundary is obtained by applying the “mod 2 union rule” (also known as the “odd-even rule”): A point is in the boundary of a `MultiCurve` if it is in the boundaries of an odd number of `Curve` elements.
- A `MultiCurve` is closed if all of its elements are closed.
- The boundary of a closed `MultiCurve` is always empty.

11.5.2.11 MultiLineString Class

A `MultiLineString` is a `MultiCurve` geometry collection composed of `LineString` elements.

`MultiLineString` Examples

- On a region map, a `MultiLineString` could represent a river system or a highway system.

11.5.2.12 MultiSurface Class

A `MultiSurface` is a geometry collection composed of surface elements. `MultiSurface` is a noninstantiable class. Its only instantiable subclass is `MultiPolygon`.

`MultiSurface` Assertions

- Surfaces within a `MultiSurface` have no interiors that intersect.
- Surfaces within a `MultiSurface` have boundaries that intersect at most at a finite number of points.

11.5.2.13 MultiPolygon Class

A `MultiPolygon` is a `MultiSurface` object composed of `Polygon` elements.

`MultiPolygon` Examples

- On a region map, a `MultiPolygon` could represent a system of lakes.

`MultiPolygon` Assertions

- A `MultiPolygon` has no two `Polygon` elements with interiors that intersect.
- A `MultiPolygon` has no two `Polygon` elements that cross (crossing is also forbidden by the previous assertion), or that touch at an infinite number of points.
- A `MultiPolygon` may not have cut lines, spikes, or punctures. A `MultiPolygon` is a regular, closed point set.
- A `MultiPolygon` that has more than one `Polygon` has an interior that is not connected. The number of connected components of the interior of a `MultiPolygon` is equal to the number of `Polygon` values in the `MultiPolygon`.

`MultiPolygon` Properties

- A `MultiPolygon` is a two-dimensional geometry.
- A `MultiPolygon` boundary is a set of closed curves (`LineString` values) corresponding to the boundaries of its `Polygon` elements.
- Each `Curve` in the boundary of the `MultiPolygon` is in the boundary of exactly one `Polygon` element.
- Every `Curve` in the boundary of an `Polygon` element is in the boundary of the `MultiPolygon`.

11.5.3 Supported Spatial Data Formats

Two standard spatial data formats are used to represent geometry objects in queries:

- Well-Known Text (WKT) format
- Well-Known Binary (WKB) format

Internally, MySQL stores geometry values in a format that is not identical to either WKT or WKB format. (Internal format is like WKB but with an initial 4 bytes to indicate the SRID.)

There are functions available to convert between different data formats; see Section 12.15.6, “Geometry Format Conversion Functions”.

The following sections describe the spatial data formats MySQL uses:

- Well-Known Text (WKT) Format
- Well-Known Binary (WKB) Format
- Internal Geometry Storage Format

Well-Known Text (WKT) Format

The Well-Known Text (WKT) representation of geometry values is designed for exchanging geometry data in ASCII form. The OpenGIS specification provides a Backus-Naur grammar that specifies the formal production rules for writing WKT values (see Section 11.5, “Extensions for Spatial Data”).

Examples of WKT representations of geometry objects:

- A `Point`:

```
POINT(15 20)
```

The point coordinates are specified with no separating comma. This differs from the syntax for the SQL `Point()` function, which requires a comma between the coordinates. Take care to use the syntax appropriate to the context of a given spatial operation. For example, the following statements both use `ST_X()` to extract the X-coordinate from a `Point` object. The first produces the object

directly using the `Point()` function. The second uses a WKT representation converted to a `Point` with `ST_GeomFromText()`.

```
mysql> SELECT ST_X(Point(15, 20));
+--------------------+
| ST_X(POINT(15, 20)) |
+--------------------+
|                 15 |
+--------------------+

mysql> SELECT ST_X(ST_GeomFromText('POINT(15 20)'));
+--------------------------------------+
| ST_X(ST_GeomFromText('POINT(15 20)')) |
+--------------------------------------+
|                                   15 |
+--------------------------------------+
```

- A `LineString` with four points:

```
LINESTRING(0 0, 10 10, 20 25, 50 60)
```

The point coordinate pairs are separated by commas.

- A `Polygon` with one exterior ring and one interior ring:

```
POLYGON((0 0,10 0,10 10,0 10,0 0),(5 5,7 5,7 7,5 7, 5 5))
```

- A `MultiPoint` with three `Point` values:

```
MULTIPOINT(0 0, 20 20, 60 60)
```

Spatial functions such as `ST_MPointFromText()` and `ST_GeomFromText()` that accept WKT-format representations of `MultiPoint` values permit individual points within values to be surrounded by parentheses. For example, both of the following function calls are valid:

```
ST_MPointFromText('MULTIPOINT (1 1, 2 2, 3 3)')
ST_MPointFromText('MULTIPOINT ((1 1), (2 2), (3 3))')
```

- A `MultiLineString` with two `LineString` values:

```
MULTILINESTRING((10 10, 20 20), (15 15, 30 15))
```

- A `MultiPolygon` with two `Polygon` values:

```
MULTIPOLYGON(((0 0,10 0,10 10,0 10,0 0)),((5 5,7 5,7 7,5 7, 5 5)))
```

- A `GeometryCollection` consisting of two `Point` values and one `LineString`:

```
GEOMETRYCOLLECTION(POINT(10 10), POINT(30 30), LINESTRING(15 15, 20 20))
```

Well-Known Binary (WKB) Format

The Well-Known Binary (WKB) representation of geometric values is used for exchanging geometry data as binary streams represented by `BLOB` values containing geometric WKB information. This format is defined by the OpenGIS specification (see Section 11.5, “Extensions for Spatial Data”). It is also defined in the ISO *SQL/MM Part 3: Spatial* standard.

WKB uses 1-byte unsigned integers, 4-byte unsigned integers, and 8-byte double-precision numbers (IEEE 754 format). A byte is eight bits.

For example, a WKB value that corresponds to `POINT(1 -1)` consists of this sequence of 21 bytes, each represented by two hexadecimal digits:

```
0101000000000000000000F03F000000000000F0BF
```

The sequence consists of the components shown in the following table.

Table 11.1 WKB Components Example

Component	Size	Value
Byte order	1 byte	`01`
WKB type	4 bytes	`01000000`
X coordinate	8 bytes	`000000000000F03F`
Y coordinate	8 bytes	`000000000000F0BF`

Component representation is as follows:

- The byte order indicator is either 1 or 0 to signify little-endian or big-endian storage. The little-endian and big-endian byte orders are also known as Network Data Representation (NDR) and External Data Representation (XDR), respectively.
- The WKB type is a code that indicates the geometry type. MySQL uses values from 1 through 7 to indicate `Point`, `LineString`, `Polygon`, `MultiPoint`, `MultiLineString`, `MultiPolygon`, and `GeometryCollection`.
- A `Point` value has X and Y coordinates, each represented as a double-precision value.

WKB values for more complex geometry values have more complex data structures, as detailed in the OpenGIS specification.

Internal Geometry Storage Format

MySQL stores geometry values using 4 bytes to indicate the SRID followed by the WKB representation of the value. For a description of WKB format, see Well-Known Binary (WKB) Format.

For the WKB part, these MySQL-specific considerations apply:

- The byte-order indicator byte is 1 because MySQL stores geometries as little-ending values.
- MySQL supports geometry types of `Point`, `LineString`, `Polygon`, `MultiPoint`, `MultiLineString`, `MultiPolygon`, and `GeometryCollection`. Other geometry types are not supported.
- Only `GeometryCollection` can be empty. Such a value is stored with 0 elements.
- Polygon rings can be specified both clockwise and counterclockwise. MySQL flips the rings automatically when reading data.

Cartesian coordinates are stored in the length unit of the spatial reference system, with X values in the X coordinates and Y values in the Y coordinates. Axis directions are those specified by the spatial reference system.

Geographic coordinates are stored in the angle unit of the spatial reference system, with longitudes in the X coordinates and latitudes in the Y coordinates. Axis directions and the meridian are those specified by the spatial reference system.

The `LENGTH()` function returns the space in bytes required for value storage. Example:

```
mysql> SET @g = ST_GeomFromText('POINT(1 -1)');
mysql> SELECT LENGTH(@g);
+------------+
| LENGTH(@g) |
+------------+
|         25 |
+------------+
mysql> SELECT HEX(@g);
+----------------------------------------------------+
| HEX(@g)                                            |
+----------------------------------------------------+
| 000000000101000000000000000000F03F000000000000F0BF |
+----------------------------------------------------+
```

The value length is 25 bytes, made up of these components (as can be seen from the hexadecimal value):

- 4 bytes for integer SRID (0)
- 1 byte for integer byte order (1 = little-endian)
- 4 bytes for integer type information (1 = `Point`)
- 8 bytes for double-precision X coordinate (1)
- 8 bytes for double-precision Y coordinate (−1)

11.5.4 Geometry Well-Formedness and Validity

For geometry values, MySQL distinguishes between the concepts of syntactically well-formed and geometrically valid.

A geometry is syntactically well-formed if it satisfies conditions such as those in this (nonexhaustive) list:

- Linestrings have at least two points
- Polygons have at least one ring
- Polygon rings are closed (first and last points the same)
- Polygon rings have at least 4 points (minimum polygon is a triangle with first and last points the same)
- Collections are not empty (except `GeometryCollection`)

A geometry is geometrically valid if it is syntactically well-formed and satisfies conditions such as those in this (nonexhaustive) list:

- Polygons are not self-intersecting
- Polygon interior rings are inside the exterior ring
- Multipolygons do not have overlapping polygons

Spatial functions fail if a geometry is not syntactically well-formed. Spatial import functions that parse WKT or WKB values raise an error for attempts to create a geometry that is not syntactically well-formed. Syntactic well-formedness is also checked for attempts to store geometries into tables.

It is permitted to insert, select, and update geometrically invalid geometries, but they must be syntactically well-formed. Due to the computational expense, MySQL does not check explicitly for geometric validity. Spatial computations may detect some cases of invalid geometries and raise an error, but they may also return an undefined result without detecting the invalidity. Applications that require geometically valid geometries should check them using the `ST_IsValid()` function.

11.5.5 Creating Spatial Columns

MySQL provides a standard way of creating spatial columns for geometry types, for example, with `CREATE TABLE` or `ALTER TABLE`. Spatial columns are supported for `MyISAM`, `InnoDB`, `NDB`, and `ARCHIVE` tables. See also the notes about spatial indexes under Section 11.5.9, “Creating Spatial Indexes”.

- Use the `CREATE TABLE` statement to create a table with a spatial column:

```
CREATE TABLE geom (g GEOMETRY);
```

- Use the `ALTER TABLE` statement to add or drop a spatial column to or from an existing table:

```
ALTER TABLE geom ADD pt POINT;
ALTER TABLE geom DROP pt;
```

11.5.6 Populating Spatial Columns

After you have created spatial columns, you can populate them with spatial data.

Values should be stored in internal geometry format, but you can convert them to that format from either Well-Known Text (WKT) or Well-Known Binary (WKB) format. The following examples demonstrate how to insert geometry values into a table by converting WKT values to internal geometry format:

- Perform the conversion directly in the `INSERT` statement:

```
INSERT INTO geom VALUES (ST_GeomFromText('POINT(1 1)'));

SET @g = 'POINT(1 1)';
INSERT INTO geom VALUES (ST_GeomFromText(@g));
```

- Perform the conversion prior to the `INSERT`:

```
SET @g = ST_GeomFromText('POINT(1 1)');
INSERT INTO geom VALUES (@g);
```

The following examples insert more complex geometries into the table:

```
SET @g = 'LINESTRING(0 0,1 1,2 2)';
INSERT INTO geom VALUES (ST_GeomFromText(@g));

SET @g = 'POLYGON((0 0,10 0,10 10,0 10,0 0),(5 5,7 5,7 7,5 7, 5 5))';
INSERT INTO geom VALUES (ST_GeomFromText(@g));

SET @g =
'GEOMETRYCOLLECTION(POINT(1 1),LINESTRING(0 0,1 1,2 2,3 3,4 4))';
INSERT INTO geom VALUES (ST_GeomFromText(@g));
```

The preceding examples use `ST_GeomFromText()` to create geometry values. You can also use type-specific functions:

```
SET @g = 'POINT(1 1)';
INSERT INTO geom VALUES (ST_PointFromText(@g));

SET @g = 'LINESTRING(0 0,1 1,2 2)';
INSERT INTO geom VALUES (ST_LineStringFromText(@g));

SET @g = 'POLYGON((0 0,10 0,10 10,0 10,0 0),(5 5,7 5,7 7,5 7, 5 5))';
INSERT INTO geom VALUES (ST_PolygonFromText(@g));
```

```
SET @g =
'GEOMETRYCOLLECTION(POINT(1 1),LINESTRING(0 0,1 1,2 2,3 3,4 4))';
INSERT INTO geom VALUES (ST_GeomCollFromText(@g));
```

A client application program that wants to use WKB representations of geometry values is responsible for sending correctly formed WKB in queries to the server. There are several ways to satisfy this requirement. For example:

- Inserting a `POINT(1 1)` value with hex literal syntax:

```
mysql> INSERT INTO geom VALUES
       (ST_GeomFromWKB(X'0101000000000000000000F03F000000000000F03F'));
```

- An ODBC application can send a WKB representation, binding it to a placeholder using an argument of `BLOB` type:

```
INSERT INTO geom VALUES (ST_GeomFromWKB(?))
```

  Other programming interfaces may support a similar placeholder mechanism.

- In a C program, you can escape a binary value using `mysql_real_escape_string_quote()` and include the result in a query string that is sent to the server. See Section 27.7.7.56, “mysql_real_escape_string_quote()”.

11.5.7 Fetching Spatial Data

Geometry values stored in a table can be fetched in internal format. You can also convert them to WKT or WKB format.

- Fetching spatial data in internal format:

  Fetching geometry values using internal format can be useful in table-to-table transfers:

```
CREATE TABLE geom2 (g GEOMETRY) SELECT g FROM geom;
```

- Fetching spatial data in WKT format:

  The `ST_AsText()` function converts a geometry from internal format to a WKT string.

```
SELECT ST_AsText(g) FROM geom;
```

- Fetching spatial data in WKB format:

  The `ST_AsBinary()` function converts a geometry from internal format to a `BLOB` containing the WKB value.

```
SELECT ST_AsBinary(g) FROM geom;
```

11.5.8 Optimizing Spatial Analysis

For `MyISAM` and `InnoDB` tables, search operations in columns containing spatial data can be optimized using `SPATIAL` indexes. The most typical operations are:

- Point queries that search for all objects that contain a given point
- Region queries that search for all objects that overlap a given region

MySQL uses **R-Trees with quadratic splitting** for `SPATIAL` indexes on spatial columns. A `SPATIAL` index is built using the minimum bounding rectangle (MBR) of a geometry. For most geometries, the

MBR is a minimum rectangle that surrounds the geometries. For a horizontal or a vertical linestring, the MBR is a rectangle degenerated into the linestring. For a point, the MBR is a rectangle degenerated into the point.

It is also possible to create normal indexes on spatial columns. In a non-`SPATIAL` index, you must declare a prefix for any spatial column except for `POINT` columns.

`MyISAM` and `InnoDB` support both `SPATIAL` and non-`SPATIAL` indexes. Other storage engines support non-`SPATIAL` indexes, as described in Section 13.1.12, “CREATE INDEX Syntax”.

11.5.9 Creating Spatial Indexes

For `MyISAM` and `InnoDB` tables, MySQL can create spatial indexes using syntax similar to that for creating regular indexes, but using the `SPATIAL` keyword. Columns in spatial indexes must be declared `NOT NULL`. The following examples demonstrate how to create spatial indexes:

- With `CREATE TABLE`:

```
CREATE TABLE geom (g GEOMETRY NOT NULL, SPATIAL INDEX(g)) ENGINE=MyISAM;
```

- With `ALTER TABLE`:

```
ALTER TABLE geom ADD SPATIAL INDEX(g);
```

- With `CREATE INDEX`:

```
CREATE SPATIAL INDEX sp_index ON geom (g);
```

`SPATIAL INDEX` creates an R-tree index. For storage engines that support nonspatial indexing of spatial columns, the engine creates a B-tree index. A B-tree index on spatial values is useful for exact-value lookups, but not for range scans.

For more information on indexing spatial columns, see Section 13.1.12, “CREATE INDEX Syntax”.

To drop spatial indexes, use `ALTER TABLE` or `DROP INDEX`:

- With `ALTER TABLE`:

```
ALTER TABLE geom DROP INDEX g;
```

- With `DROP INDEX`:

```
DROP INDEX sp_index ON geom;
```

Example: Suppose that a table `geom` contains more than 32,000 geometries, which are stored in the column `g` of type `GEOMETRY`. The table also has an `AUTO_INCREMENT` column `fid` for storing object ID values.

```
mysql> DESCRIBE geom;
+-------+----------+------+-----+---------+----------------+
| Field | Type     | Null | Key | Default | Extra          |
+-------+----------+------+-----+---------+----------------+
| fid   | int(11)  |      | PRI | NULL    | auto_increment |
| g     | geometry |      |     |         |                |
+-------+----------+------+-----+---------+----------------+
2 rows in set (0.00 sec)

mysql> SELECT COUNT(*) FROM geom;
+----------+
```

```
| count(*) |
+----------+
|    32376 |
+----------+
1 row in set (0.00 sec)
```

To add a spatial index on the column `g`, use this statement:

```
mysql> ALTER TABLE geom ADD SPATIAL INDEX(g) ENGINE=MyISAM;
Query OK, 32376 rows affected (4.05 sec)
Records: 32376  Duplicates: 0  Warnings: 0
```

11.5.10 Using Spatial Indexes

The optimizer investigates whether available spatial indexes can be involved in the search for queries that use a function such as `MBRContains()` or `MBRWithin()` in the `WHERE` clause. The following query finds all objects that are in the given rectangle:

```
mysql> SET @poly =
    -> 'Polygon((30000 15000,
                 31000 15000,
                 31000 16000,
                 30000 16000,
                 30000 15000))';
mysql> SELECT fid,ST_AsText(g) FROM geom WHERE
    -> MBRContains(ST_GeomFromText(@poly),g);
+-----+-------------------------------------------------------------+
| fid | ST_AsText(g)                                                |
+-----+-------------------------------------------------------------+
|  21 | LINESTRING(30350.4 15828.8,30350.6 15845,30333.8 15845,30 ... |
|  22 | LINESTRING(30350.6 15871.4,30350.6 15887.8,30334 15887.8, ... |
|  23 | LINESTRING(30350.6 15914.2,30350.6 15930.4,30334 15930.4, ... |
|  24 | LINESTRING(30290.2 15823,30290.2 15839.4,30273.4 15839.4, ... |
|  25 | LINESTRING(30291.4 15866.2,30291.6 15882.4,30274.8 15882. ... |
|  26 | LINESTRING(30291.6 15918.2,30291.6 15934.4,30275 15934.4, ... |
| 249 | LINESTRING(30337.8 15938.6,30337.8 15946.8,30320.4 15946. ... |
|   1 | LINESTRING(30250.4 15129.2,30248.8 15138.4,30238.2 15136. ... |
|   2 | LINESTRING(30220.2 15122.8,30217.2 15137.8,30207.6 15136, ... |
|   3 | LINESTRING(30179 15114.4,30176.6 15129.4,30167 15128,3016 ... |
|   4 | LINESTRING(30155.2 15121.4,30140.4 15118.6,30142 15109,30 ... |
|   5 | LINESTRING(30192.4 15085,30177.6 15082.2,30179.2 15072.4, ... |
|   6 | LINESTRING(30244 15087,30229 15086.2,30229.4 15076.4,3024 ... |
|   7 | LINESTRING(30200.6 15059.4,30185.6 15058.6,30186 15048.8, ... |
|  10 | LINESTRING(30179.6 15017.8,30181 15002.8,30190.8 15003.6, ... |
|  11 | LINESTRING(30154.2 15000.4,30168.6 15004.8,30166 15014.2, ... |
|  13 | LINESTRING(30105 15065.8,30108.4 15050.8,30118 15053,3011 ... |
| 154 | LINESTRING(30276.2 15143.8,30261.4 15141,30263 15131.4,30 ... |
| 155 | LINESTRING(30269.8 15084,30269.4 15093.4,30258.6 15093,30 ... |
| 157 | LINESTRING(30128.2 15011,30113.2 15010.2,30113.6 15000.4, ... |
+-----+-------------------------------------------------------------+
20 rows in set (0.00 sec)
```

Use `EXPLAIN` to check the way this query is executed:

```
mysql> SET @poly =
    -> 'Polygon((30000 15000,
                 31000 15000,
                 31000 16000,
                 30000 16000,
                 30000 15000))';
mysql> EXPLAIN SELECT fid,ST_AsText(g) FROM geom WHERE
    -> MBRContains(ST_GeomFromText(@poly),g)\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: geom
         type: range
```

```
possible_keys: g
          key: g
      key_len: 32
          ref: NULL
         rows: 50
        Extra: Using where
1 row in set (0.00 sec)
```

Check what would happen without a spatial index:

```
mysql> SET @poly =
    -> 'Polygon((30000 15000,
                 31000 15000,
                 31000 16000,
                 30000 16000,
                 30000 15000))';
mysql> EXPLAIN SELECT fid,ST_AsText(g) FROM g IGNORE INDEX (g) WHERE
    -> MBRContains(ST_GeomFromText(@poly),g)\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: geom
         type: ALL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 32376
        Extra: Using where
1 row in set (0.00 sec)
```

Executing the `SELECT` statement without the spatial index yields the same result but causes the execution time to rise from 0.00 seconds to 0.46 seconds:

```
mysql> SET @poly =
    -> 'Polygon((30000 15000,
                 31000 15000,
                 31000 16000,
                 30000 16000,
                 30000 15000))';
mysql> SELECT fid,ST_AsText(g) FROM geom IGNORE INDEX (g) WHERE
    -> MBRContains(ST_GeomFromText(@poly),g);
+-----+-------------------------------------------------------------+
| fid | ST_AsText(g)                                                |
+-----+-------------------------------------------------------------+
|   1 | LINESTRING(30250.4 15129.2,30248.8 15138.4,30238.2 15136. ... |
|   2 | LINESTRING(30220.2 15122.8,30217.2 15137.8,30207.6 15136, ... |
|   3 | LINESTRING(30179 15114.4,30176.6 15129.4,30167 15128,3016 ... |
|   4 | LINESTRING(30155.2 15121.4,30140.4 15118.6,30142 15109,30 ... |
|   5 | LINESTRING(30192.4 15085,30177.6 15082.2,30179.2 15072.4, ... |
|   6 | LINESTRING(30244 15087,30229 15086.2,30229.4 15076.4,3024 ... |
|   7 | LINESTRING(30200.6 15059.4,30185.6 15058.6,30186 15048.8, ... |
|  10 | LINESTRING(30179.6 15017.8,30181 15002.8,30190.8 15003.6, ... |
|  11 | LINESTRING(30154.2 15000.4,30168.6 15004.8,30166 15014.2, ... |
|  13 | LINESTRING(30105 15065.8,30108.4 15050.8,30118 15053,3011 ... |
|  21 | LINESTRING(30350.4 15828.8,30350.6 15845,30333.8 15845,30 ... |
|  22 | LINESTRING(30350.6 15871.4,30350.6 15887.8,30334 15887.8, ... |
|  23 | LINESTRING(30350.6 15914.2,30350.6 15930.4,30334 15930.4, ... |
|  24 | LINESTRING(30290.2 15823,30290.2 15839.4,30273.4 15839.4, ... |
|  25 | LINESTRING(30291.4 15866.2,30291.6 15882.4,30274.8 15882. ... |
|  26 | LINESTRING(30291.6 15918.2,30291.6 15934.4,30275 15934.4, ... |
| 154 | LINESTRING(30276.2 15143.8,30261.4 15141,30263 15131.4,30 ... |
| 155 | LINESTRING(30269.8 15084,30269.4 15093.4,30258.6 15093,30 ... |
| 157 | LINESTRING(30128.2 15011,30113.2 15010.2,30113.6 15000.4, ... |
| 249 | LINESTRING(30337.8 15938.6,30337.8 15946.8,30320.4 15946. ... |
+-----+-------------------------------------------------------------+
20 rows in set (0.46 sec)
```

11.6 The JSON Data Type

- Creating JSON Values
- Normalization, Merging, and Autowrapping of JSON Values
- Searching and Modifying JSON Values
- Comparison and Ordering of JSON Values
- Converting between JSON and non-JSON values
- Aggregation of JSON Values

MySQL supports a native `JSON` data type defined by RFC 7159 that enables efficient access to data in JSON (JavaScript Object Notation) documents. The `JSON` data type provides these advantages over storing JSON-format strings in a string column:

- Automatic validation of JSON documents stored in `JSON` columns. Invalid documents produce an error.
- Optimized storage format. JSON documents stored in `JSON` columns are converted to an internal format that permits quick read access to document elements. When the server later must read a JSON value stored in this binary format, the value need not be parsed from a text representation. The binary format is structured to enable the server to look up subobjects or nested values directly by key or array index without reading all values before or after them in the document.

Note

This discussion uses `JSON` in monotype to indicate specifically the JSON data type and “JSON” in regular font to indicate JSON data in general.

The space required to store a `JSON` document is roughly the same as for `LONGBLOB` or `LONGTEXT`; see Section 11.8, “Data Type Storage Requirements”, for more information. It is important to keep in mind that the size of any JSON document stored in a `JSON` column is limited to the value of the `max_allowed_packet` system variable. (When the server is manipulating a JSON value internally in memory, it can be larger than this; the limit applies when the server stores it.) You can obtain the amount of space required to store a JSON document using the `JSON_STORAGE_SIZE()` function (introduced in MySQL 8.0.2); note that for a `JSON` column, the storage size—and thus the value returned by this function—is that used by the column prior to any partial updates that may have been performed on it (see the discussion of the JSON partial update optimization later in this section).

A `JSON` column cannot have a default value.

Along with the `JSON` data type, a set of SQL functions is available to enable operations on JSON values, such as creation, manipulation, and searching. The following discussion shows examples of these operations. For details about individual functions, see Section 12.16, “JSON Functions”.

A set of spatial functions for operating on GeoJSON values is also available. See Section 12.15.11, “Spatial GeoJSON Functions”.

`JSON` columns, like columns of other binary types, are not indexed directly; instead, you can create an index on a generated column that extracts a scalar value from the `JSON` column. See Indexing a Generated Column to Provide a JSON Column Index, for a detailed example.

The MySQL optimizer also looks for compatible indexes on virtual columns that match JSON expressions.

In MySQL 8.0 (8.0.2 and later), the optimizer can perform a partial, in-place update of a `JSON` column instead of removing the old document and writing the new document in its entirety to the column. This optimization can be performed for an update that meets the following conditions:

- The column being updated was declared as `JSON`.

- The `UPDATE` statement uses any of the three functions `JSON_SET()`, `JSON_REPLACE()`, or `JSON_REMOVE()` to update the column. A direct assignment of the column value (for example, `UPDATE mytable SET jcol = '{"a": 10, "b": 25'}`) cannot be performed as a partial update.

  Updates of multiple `JSON` columns in a single `UPDATE` statement can be optimized in this fashion; MySQL can perform partial updates of only those columns whose values are updated using the three functions just listed.

- The input column and the target column must be the same column; a statement such as `UPDATE mytable SET jcol1 = JSON_SET(jcol2, '$.a', 100)` cannot be performed as a partial update.

  The update can use nested calls to any of the functions listed in the previous item, in any combination, as long as the input and target columns are the same.

- All changes replace existing array or object values with new ones, and do not add any new elements to the parent object or array.

- The value being replaced must be at least as large as the replacement value. In other words, the new value cannot be any larger than the old one.

  A possible exception to this requirement jus stated can occur when a previous partial update has left sufficient space for the larger value. You can use the function `JSON_STORAGE_FREE()` see how much space was freed by the most recent partial update of a `JSON` column.

The next few sections provide basic information regarding the creation and manipulation of JSON values.

Creating JSON Values

A JSON array contains a list of values separated by commas and enclosed within `[` and `]` characters:

```
["abc", 10, null, true, false]
```

A JSON object contains a set of key-value pairs separated by commas and enclosed within `{` and `}` characters:

```
{"k1": "value", "k2": 10}
```

As the examples illustrate, JSON arrays and objects can contain scalar values that are strings or numbers, the JSON null literal, or the JSON boolean true or false literals. Keys in JSON objects must be strings. Temporal (date, time, or datetime) scalar values are also permitted:

```
["12:18:29.000000", "2015-07-29", "2015-07-29 12:18:29.000000"]
```

Nesting is permitted within JSON array elements and JSON object key values:

```
[99, {"id": "HK500", "cost": 75.99}, ["hot", "cold"]]
{"k1": "value", "k2": [10, 20]}
```

You can also obtain JSON values from a number of functions supplied by MySQL for this purpose (see Section 12.16.2, “Functions That Create JSON Values”) as well as by casting values of other types to the `JSON` type using `CAST(value AS JSON)` (see Converting between JSON and non-JSON values). The next several paragraphs describe how MySQL handles JSON values provided as input.

In MySQL, JSON values are written as strings. MySQL parses any string used in a context that requires a JSON value, and produces an error if it is not valid as JSON. These contexts include inserting a value into a column that has the `JSON` data type and passing an argument to a function that

expects a JSON value (usually shown as *`json_doc`* or *`json_val`* in the documentation for MySQL JSON functions), as the following examples demonstrate:

- Attempting to insert a value into a `JSON` column succeeds if the value is a valid JSON value, but fails if it is not:

```
mysql> CREATE TABLE t1 (jdoc JSON);
Query OK, 0 rows affected (0.20 sec)

mysql> INSERT INTO t1 VALUES('{"key1": "value1", "key2": "value2"}');
Query OK, 1 row affected (0.01 sec)

mysql> INSERT INTO t1 VALUES('[1, 2,');
ERROR 3140 (22032) at line 2: Invalid JSON text:
"Invalid value." at position 6 in value (or column) '[1, 2,'.
```

  Positions for “at position *`N`*” in such error messages are 0-based, but should be considered rough indications of where the problem in a value actually occurs.

- The `JSON_TYPE()` function expects a JSON argument and attempts to parse it into a JSON value. It returns the value's JSON type if it is valid and produces an error otherwise:

```
mysql> SELECT JSON_TYPE('["a", "b", 1]');
+---------------------------+
| JSON_TYPE('["a", "b", 1]') |
+---------------------------+
| ARRAY                     |
+---------------------------+

mysql> SELECT JSON_TYPE('"hello"');
+----------------------+
| JSON_TYPE('"hello"') |
+----------------------+
| STRING               |
+----------------------+

mysql> SELECT JSON_TYPE('hello');
ERROR 3146 (22032): Invalid data type for JSON data in argument 1
to function json_type; a JSON string or JSON type is required.
```

MySQL handles strings used in JSON context using the `utf8mb4` character set and `utf8mb4_bin` collation. Strings in other character sets are converted to `utf8mb4` as necessary. (For strings in the `ascii` or `utf8` character sets, no conversion is needed because `ascii` and `utf8` are subsets of `utf8mb4`.)

As an alternative to writing JSON values using literal strings, functions exist for composing JSON values from component elements. `JSON_ARRAY()` takes a (possibly empty) list of values and returns a JSON array containing those values:

```
mysql> SELECT JSON_ARRAY('a', 1, NOW());
+---------------------------------------+
| JSON_ARRAY('a', 1, NOW())             |
+---------------------------------------+
| ["a", 1, "2015-07-27 09:43:47.000000"] |
+---------------------------------------+
```

`JSON_OBJECT()` takes a (possibly empty) list of key-value pairs and returns a JSON object containing those pairs:

```
mysql> SELECT JSON_OBJECT('key1', 1, 'key2', 'abc');
+--------------------------------------+
| JSON_OBJECT('key1', 1, 'key2', 'abc') |
+--------------------------------------+
| {"key1": 1, "key2": "abc"}           |
```

```
+--------------------------------------+
```

`JSON_MERGE()` takes two or more JSON documents and returns the combined result:

```
mysql> SELECT JSON_MERGE('["a", 1]', '{"key": "value"}');
+-----------------------------------------+
| JSON_MERGE('["a", 1]', '{"key": "value"}') |
+-----------------------------------------+
| ["a", 1, {"key": "value"}]              |
+-----------------------------------------+
```

For information about the merging rules, see Normalization, Merging, and Autowrapping of JSON Values.

JSON values can be assigned to user-defined variables:

```
mysql> SET @j = JSON_OBJECT('key', 'value');
mysql> SELECT @j;
+------------------+
| @j               |
+------------------+
| {"key": "value"} |
+------------------+
```

However, user-defined variables cannot be of `JSON` data type, so although `@j` in the preceding example looks like a JSON value and has the same character set and collation as a JSON value, it does *not* have the `JSON` data type. Instead, the result from `JSON_OBJECT()` is converted to a string when assigned to the variable.

Strings produced by converting JSON values have a character set of `utf8mb4` and a collation of `utf8mb4_bin`:

```
mysql> SELECT CHARSET(@j), COLLATION(@j);
+-------------+---------------+
| CHARSET(@j) | COLLATION(@j) |
+-------------+---------------+
| utf8mb4     | utf8mb4_bin   |
+-------------+---------------+
```

Because `utf8mb4_bin` is a binary collation, comparison of JSON values is case sensitive.

```
mysql> SELECT JSON_ARRAY('x') = JSON_ARRAY('X');
+-----------------------------------+
| JSON_ARRAY('x') = JSON_ARRAY('X') |
+-----------------------------------+
|                                 0 |
+-----------------------------------+
```

Case sensitivity also applies to the JSON `null`, `true`, and `false` literals, which always must be written in lowercase:

```
mysql> SELECT JSON_VALID('null'), JSON_VALID('Null'), JSON_VALID('NULL');
+--------------------+--------------------+--------------------+
| JSON_VALID('null') | JSON_VALID('Null') | JSON_VALID('NULL') |
+--------------------+--------------------+--------------------+
|                  1 |                  0 |                  0 |
+--------------------+--------------------+--------------------+

mysql> SELECT CAST('null' AS JSON);
+----------------------+
| CAST('null' AS JSON) |
+----------------------+
| null                 |
+----------------------+
```

```
1 row in set (0.00 sec)

mysql> SELECT CAST('NULL' AS JSON);
ERROR 3141 (22032): Invalid JSON text in argument 1 to function cast_as_json:
"Invalid value." at position 0 in 'NULL'.
```

Case sensitivity of the JSON literals differs from that of the SQL `NULL`, `TRUE`, and `FALSE` literals, which can be written in any lettercase:

```
mysql> SELECT ISNULL(null), ISNULL(Null), ISNULL(NULL);
+--------------+--------------+--------------+
| ISNULL(null) | ISNULL(Null) | ISNULL(NULL) |
+--------------+--------------+--------------+
|            1 |            1 |            1 |
+--------------+--------------+--------------+
```

Sometimes it may be necessary or desirable to insert quote characters (`"` or `'`) into a JSON document. Assume for this example that you want to insert some JSON objects containing strings representing sentences that state some facts about MySQL, each paired with an appropriate keyword, into a table created using the SQL statement shown here:

```
mysql> CREATE TABLE facts (sentence JSON);
```

Among these keyword-sentence pairs is this one:

```
mascot: The MySQL mascot is a dolphin named "Sakila".
```

One way to insert this as a JSON object into the `facts` table is to use the MySQL `JSON_OBJECT()` function. In this case, you must escape each quote character using a backslash, as shown here:

```
mysql> INSERT INTO facts VALUES
    >   (JSON_OBJECT("mascot", "Our mascot is a dolphin named \"Sakila\"."));
```

This does not work in the same way if you insert the value as a JSON object literal, in which case, you must use the double backslash escape sequence, like this:

```
mysql> INSERT INTO facts VALUES
    >   ('{"mascot": "Our mascot is a dolphin named \\"Sakila\\"."}');
```

Using the double backslash keeps MySQL from performing escape sequence processing, and instead causes it to pass the string literal to the storage engine for processing. After inserting the JSON object in either of the ways just shown, you can see that the backslashes are present in the JSON column value by doing a simple `SELECT`, like this:

```
mysql> SELECT sentence FROM facts;
+-----------------------------------------------------------+
| sentence                                                  |
+-----------------------------------------------------------+
| {"mascot": "Our mascot is a dolphin named \"Sakila\"."} |
+-----------------------------------------------------------+
```

To look up this particular sentence employing `mascot` as the key, you can use the column-path operator `->`, as shown here:

```
mysql> SELECT col->"$.mascot" FROM qtest;
+-----------------------------------------+
| col->"$.mascot"                         |
+-----------------------------------------+
| "Our mascot is a dolphin named \"Sakila\"." |
+-----------------------------------------+
1 row in set (0.00 sec)
```

This leaves the backslashes intact, along with the surrounding quote marks. To display the desired value using `mascot` as the key, but without including the surrounding quote marks or any escapes, use the inline path operator `->>`, like this:

```
mysql> SELECT sentence->>"$.mascot" FROM facts;
+---------------------------------------+
| sentence->>"$.mascot"                 |
+---------------------------------------+
| Our mascot is a dolphin named "Sakila". |
+---------------------------------------+
```

Note

The previous example does not work as shown if the `NO_BACKSLASH_ESCAPES` server SQL mode is enabled. If this mode is set, a single backslash instead of double backslashes can be used to insert the JSON object literal, and the backslashes are preserved. If you use the `JSON_OBJECT()` function when performing the insert and this mode is set, you must alternate single and double quotes, like this:

```
mysql> INSERT INTO facts VALUES
    > (JSON_OBJECT('mascot', 'Our mascot is a dolphin named "Sakila".'));
```

See the description of the `JSON_UNQUOTE()` function for more information about the effects of this mode on escaped characters in JSON values.

Normalization, Merging, and Autowrapping of JSON Values

When a string is parsed and found to be a valid JSON document, it is also normalized: Members with keys that duplicate a key found earlier in the document are discarded (even if the values differ). The object value produced by the following `JSON_OBJECT()` call does not include the second `key1` element because that key name occurs earlier in the value:

```
mysql> SELECT JSON_OBJECT('key1', 1, 'key2', 'abc', 'key1', 'def');
+-----------------------------------------------------+
| JSON_OBJECT('key1', 1, 'key2', 'abc', 'key1', 'def') |
+-----------------------------------------------------+
| {"key1": 1, "key2": "abc"}                          |
+-----------------------------------------------------+
```

The normalization performed by MySQL also sorts the keys of a JSON object (for the purpose of making lookups more efficient). The result of this ordering is subject to change and not guaranteed to be consistent across releases. In addition, extra whitespace between keys, values, or elements in the original document is discarded.

MySQL functions that produce JSON values (see Section 12.16.2, “Functions That Create JSON Values”) always return normalized values.

In contexts that combine multiple arrays, the arrays are merged into a single array by concatenating arrays named later to the end of the first array. In the following example, `JSON_MERGE()` merges its arguments into a single array:

```
mysql> SELECT JSON_MERGE('[1, 2]', '["a", "b"]', '[true, false]');
+-----------------------------------------------------+
| JSON_MERGE('[1, 2]', '["a", "b"]', '[true, false]') |
+-----------------------------------------------------+
| [1, 2, "a", "b", true, false]                       |
+-----------------------------------------------------+
```

Multiple objects when merged produce a single object. If multiple objects have the same key, the value for that key in the resulting merged object is an array containing the key values:

```
mysql> SELECT JSON_MERGE('{"a": 1, "b": 2}', '{"c": 3, "a": 4}');
+--------------------------------------------------+
| JSON_MERGE('{"a": 1, "b": 2}', '{"c": 3, "a": 4}') |
+--------------------------------------------------+
| {"a": [1, 4], "b": 2, "c": 3}                    |
+--------------------------------------------------+
```

Nonarray values used in a context that requires an array value are autowrapped: The value is surrounded by `[` and `]` characters to convert it to an array. In the following statement, each argument is autowrapped as an array (`[1]`, `[2]`). These are then merged to produce a single result array:

```
mysql> SELECT JSON_MERGE('1', '2');
+----------------------+
| JSON_MERGE('1', '2') |
+----------------------+
| [1, 2]               |
+----------------------+
```

Array and object values are merged by autowrapping the object as an array and merging the two arrays:

```
mysql> SELECT JSON_MERGE('[10, 20]', '{"a": "x", "b": "y"}');
+-----------------------------------------------+
| JSON_MERGE('[10, 20]', '{"a": "x", "b": "y"}') |
+-----------------------------------------------+
| [10, 20, {"a": "x", "b": "y"}]                |
+-----------------------------------------------+
```

Searching and Modifying JSON Values

A JSON path expression selects a value within a JSON document.

Path expressions are useful with functions that extract parts of or modify a JSON document, to specify where within that document to operate. For example, the following query extracts from a JSON document the value of the member with the `name` key:

```
mysql> SELECT JSON_EXTRACT('{"id": 14, "name": "Aztalan"}', '$.name');
+---------------------------------------------------------+
| JSON_EXTRACT('{"id": 14, "name": "Aztalan"}', '$.name') |
+---------------------------------------------------------+
| "Aztalan"                                               |
+---------------------------------------------------------+
```

Path syntax uses a leading `$` character to represent the JSON document under consideration, optionally followed by selectors that indicate successively more specific parts of the document:

- A period followed by a key name names the member in an object with the given key. The key name must be specified within double quotation marks if the name without quotes is not legal within path expressions (for example, if it contains a space).
- `[N]` appended to a path that selects an array names the value at position *`N`* within the array. Array positions are integers beginning with zero.
- (MySQL 8.0.2 and later:) `[M to N]` specifies a subset or range of array values starting with the value at position *`M`*, and ending with the value at position *`N`*.

  `last` is supported as a synonym for the index of the rightmost array element. Relative addressing of array elements is also supported.
- Paths can contain `*` or `**` wildcards:
  - `.[*]` evaluates to the values of all members in a JSON object.

  - `[*]` evaluates to the values of all elements in a JSON array.
  - *`prefix**suffix`* evaluates to all paths that begin with the named prefix and end with the named suffix.
- A path that does not exist in the document (evaluates to nonexistent data) evaluates to `NULL`.

Let `$` refer to this JSON array with three elements:

```
[3, {"a": [5, 6], "b": 10}, [99, 100]]
```

Then:

- `$[0]` evaluates to `3`.
- `$[1]` evaluates to `{"a": [5, 6], "b": 10}`.
- `$[2]` evaluates to `[99, 100]`.
- `$[3]` evaluates to `NULL` (it refers to the fourth array element, which does not exist).

Because `$[1]` and `$[2]` evaluate to nonscalar values, they can be used as the basis for more-specific path expressions that select nested values. Examples:

- `$[1].a` evaluates to `[5, 6]`.
- `$[1].a[1]` evaluates to `6`.
- `$[1].b` evaluates to `10`.
- `$[2][0]` evaluates to `99`.

As mentioned previously, path components that name keys must be quoted if the unquoted key name is not legal in path expressions. Let `$` refer to this value:

```
{"a fish": "shark", "a bird": "sparrow"}
```

The keys both contain a space and must be quoted:

- `$."a fish"` evaluates to `shark`.
- `$."a bird"` evaluates to `sparrow`.

Paths that use wildcards evaluate to an array that can contain multiple values:

```
mysql> SELECT JSON_EXTRACT('{"a": 1, "b": 2, "c": [3, 4, 5]}', '$.*');
+---------------------------------------------------------+
| JSON_EXTRACT('{"a": 1, "b": 2, "c": [3, 4, 5]}', '$.*') |
+---------------------------------------------------------+
| [1, 2, [3, 4, 5]]                                       |
+---------------------------------------------------------+
mysql> SELECT JSON_EXTRACT('{"a": 1, "b": 2, "c": [3, 4, 5]}', '$.c[*]');
+------------------------------------------------------------+
| JSON_EXTRACT('{"a": 1, "b": 2, "c": [3, 4, 5]}', '$.c[*]') |
+------------------------------------------------------------+
| [3, 4, 5]                                                  |
+------------------------------------------------------------+
```

In the following example, the path `$**.b` evaluates to multiple paths (`$.a.b` and `$.c.b`) and produces an array of the matching path values:

```
mysql> SELECT JSON_EXTRACT('{"a": {"b": 1}, "c": {"b": 2}}', '$**.b');
+-----------------------------------------------------------+
| JSON_EXTRACT('{"a": {"b": 1}, "c": {"b": 2}}', '$**.b') |
+-----------------------------------------------------------+
| [1, 2]                                                    |
+-----------------------------------------------------------+
```

Beginning with MySQL 8.0.2, you can use ranges with the `to` keyword to specify subsets of JSON arrays. For example, `$[1 to 3]` includes the second, third, and fourth elements of an array, as shown here:

```
mysql> SELECT JSON_EXTRACT('[1, 2, 3, 4, 5]', '$[1 to 3]');
+---------------------------------------------+
| JSON_EXTRACT('[1, 2, 3, 4, 5]', '$[1 to 3]') |
+---------------------------------------------+
| [2, 3, 4]                                   |
+---------------------------------------------+
1 row in set (0.00 sec)
```

The syntax is *`M to N`*, where *`M`* and *`N`* are, respectively, the first and last indexes of a range of elements from a JSON array. *`N`* must be greater than *`M`*; *`M`* must be greater than or equal to 0. Array elements are indexed beginning with 0.

You can use ranges in contexts where wildcards are supported.

Beginning with MySQL 8.0.2, the `last` keyword is supported as a synonym for the index of the last element in an array. Expressions of the form `last - N` can be used for relative addressing, and within range definitions, like this:

```
mysql> SELECT JSON_EXTRACT('[1, 2, 3, 4, 5]', '$[last-3 to last-1]');
+-------------------------------------------------------+
| JSON_EXTRACT('[1, 2, 3, 4, 5]', '$[last-3 to last-1]') |
+-------------------------------------------------------+
| [2, 3, 4]                                             |
+-------------------------------------------------------+
1 row in set (0.01 sec)
```

If the path is evaluated against a value that is not an array, the result of the evaluation is the same as if the value had been wrapped in a single-element array:

```
mysql> SELECT JSON_REPLACE('"Sakila"', '$[last]', 10);
+-----------------------------------------+
| JSON_REPLACE('"Sakila"', '$[last]', 10) |
+-----------------------------------------+
| 10                                      |
+-----------------------------------------+
1 row in set (0.00 sec)
```

You can use *`column->path`* with a JSON column identifier and JSON path expression as a synonym for `JSON_EXTRACT(`*`column, path`*`)`. See Section 12.16.3, “Functions That Search JSON Values”, for more information. See also Indexing a Generated Column to Provide a JSON Column Index.

Some functions take an existing JSON document, modify it in some way, and return the resulting modified document. Path expressions indicate where in the document to make changes. For example, the `JSON_SET()`, `JSON_INSERT()`, and `JSON_REPLACE()` functions each take a JSON document, plus one or more path-value pairs that describe where to modify the document and the values to use. The functions differ in how they handle existing and nonexisting values within the document.

Consider this document:

```
mysql> SET @j = '["a", {"b": [true, false]}, [10, 20]]';
```

`JSON_SET()` replaces values for paths that exist and adds values for paths that do not exist:.

```
mysql> SELECT JSON_SET(@j, '$[1].b[0]', 1, '$[2][2]', 2);
+-----------------------------------------+
| JSON_SET(@j, '$[1].b[0]', 1, '$[2][2]', 2) |
+-----------------------------------------+
| ["a", {"b": [1, false]}, [10, 20, 2]]      |
+-----------------------------------------+
```

In this case, the path `$[1].b[0]` selects an existing value (`true`), which is replaced with the value following the path argument (`1`). The path `$[2][2]` does not exist, so the corresponding value (`2`) is added to the value selected by `$[2]`.

`JSON_INSERT()` adds new values but does not replace existing values:

```
mysql> SELECT JSON_INSERT(@j, '$[1].b[0]', 1, '$[2][2]', 2);
+-----------------------------------------+
| JSON_INSERT(@j, '$[1].b[0]', 1, '$[2][2]', 2) |
+-----------------------------------------+
| ["a", {"b": [true, false]}, [10, 20, 2]]      |
+-----------------------------------------+
```

`JSON_REPLACE()` replaces existing values and ignores new values:

```
mysql> SELECT JSON_REPLACE(@j, '$[1].b[0]', 1, '$[2][2]', 2);
+------------------------------------------+
| JSON_REPLACE(@j, '$[1].b[0]', 1, '$[2][2]', 2) |
+------------------------------------------+
| ["a", {"b": [1, false]}, [10, 20]]             |
+------------------------------------------+
```

The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

`JSON_REMOVE()` takes a JSON document and one or more paths that specify values to be removed from the document. The return value is the original document minus the values selected by paths that exist within the document:

```
mysql> SELECT JSON_REMOVE(@j, '$[2]', '$[1].b[1]', '$[1].b[1]');
+-------------------------------------------------+
| JSON_REMOVE(@j, '$[2]', '$[1].b[1]', '$[1].b[1]') |
+-------------------------------------------------+
| ["a", {"b": [true]}]                            |
+-------------------------------------------------+
```

The paths have these effects:

- `$[2]` matches `[10, 20]` and removes it.
- The first instance of `$[1].b[1]` matches `false` in the `b` element and removes it.
- The second instance of `$[1].b[1]` matches nothing: That element has already been removed, the path no longer exists, and has no effect.

Comparison and Ordering of JSON Values

JSON values can be compared using the `=`, `<`, `<=`, `>`, `>=`, `<>`, `!=`, and `<=>` operators.

The following comparison operators and functions are not yet supported with JSON values:

- `BETWEEN`
- `IN()`
- `GREATEST()`

- `LEAST()`

A workaround for the comparison operators and functions just listed is to cast JSON values to a native MySQL numeric or string data type so they have a consistent non-JSON scalar type.

Comparison of JSON values takes place at two levels. The first level of comparison is based on the JSON types of the compared values. If the types differ, the comparison result is determined solely by which type has higher precedence. If the two values have the same JSON type, a second level of comparison occurs using type-specific rules.

The following list shows the precedences of JSON types, from highest precedence to the lowest. (The type names are those returned by the `JSON_TYPE()` function.) Types shown together on a line have the same precedence. Any value having a JSON type listed earlier in the list compares greater than any value having a JSON type listed later in the list.

```
BLOB
BIT
OPAQUE
DATETIME
TIME
DATE
BOOLEAN
ARRAY
OBJECT
STRING
INTEGER, DOUBLE
NULL
```

For JSON values of the same precedence, the comparison rules are type specific:

- `BLOB`

  The first *`N`* bytes of the two values are compared, where *`N`* is the number of bytes in the shorter value. If the first *`N`* bytes of the two values are identical, the shorter value is ordered before the longer value.

- `BIT`

  Same rules as for `BLOB`.

- `OPAQUE`

  Same rules as for `BLOB`. `OPAQUE` values are values that are not classified as one of the other types.

- `DATETIME`

  A value that represents an earlier point in time is ordered before a value that represents a later point in time. If two values originally come from the MySQL `DATETIME` and `TIMESTAMP` types, respectively, they are equal if they represent the same point in time.

- `TIME`

  The smaller of two time values is ordered before the larger one.

- `DATE`

  The earlier date is ordered before the more recent date.

- `ARRAY`

  Two JSON arrays are equal if they have the same length and values in corresponding positions in the arrays are equal.

  If the arrays are not equal, their order is determined by the elements in the first position where there is a difference. The array with the smaller value in that position is ordered first. If all values of the

shorter array are equal to the corresponding values in the longer array, the shorter array is ordered first.

Example:

```
[] < ["a"] < ["ab"] < ["ab", "cd", "ef"] < ["ab", "ef"]
```

- `BOOLEAN`

  The JSON false literal is less than the JSON true literal.

- `OBJECT`

  Two JSON objects are equal if they have the same set of keys, and each key has the same value in both objects.

  Example:

  ```
  {"a": 1, "b": 2} = {"b": 2, "a": 1}
  ```

  The order of two objects that are not equal is unspecified but deterministic.

- `STRING`

  Strings are ordered lexically on the first *`N`* bytes of the `utf8mb4` representation of the two strings being compared, where *`N`* is the length of the shorter string. If the first *`N`* bytes of the two strings are identical, the shorter string is considered smaller than the longer string.

  Example:

  ```
  "a" < "ab" < "b" < "bc"
  ```

  This ordering is equivalent to the ordering of SQL strings with collation `utf8mb4_bin`. Because `utf8mb4_bin` is a binary collation, comparison of JSON values is case sensitive:

  ```
  "A" < "a"
  ```

- `INTEGER`, `DOUBLE`

  JSON values can contain exact-value numbers and approximate-value numbers. For a general discussion of these types of numbers, see Section 9.1.2, “Numeric Literals”.

  The rules for comparing native MySQL numeric types are discussed in Section 12.2, “Type Conversion in Expression Evaluation”, but the rules for comparing numbers within JSON values differ somewhat:

  - In a comparison between two columns that use the native MySQL `INT` and `DOUBLE` numeric types, respectively, it is known that all comparisons involve an integer and a double, so the integer is converted to double for all rows. That is, exact-value numbers are converted to approximate-value numbers.

  - On the other hand, if the query compares two JSON columns containing numbers, it cannot be known in advance whether numbers will be integer or double. To provide the most consistent behavior across all rows, MySQL converts approximate-value numbers to exact-value numbers. The resulting ordering is consistent and does not lose precision for the exact-value numbers. For example, given the scalars 9223372036854775805, 9223372036854775806, 9223372036854775807 and 9.223372036854776e18, the order is such as this:

    ```
    9223372036854775805 < 9223372036854775806 < 9223372036854775807
    ```

```
< 9.223372036854776e18 = 9223372036854776000 < 9223372036854776001
```

Were JSON comparisons to use the non-JSON numeric comparison rules, inconsistent ordering could occur. The usual MySQL comparison rules for numbers yield these orderings:

- Integer comparison:

```
9223372036854775805 < 9223372036854775806 < 9223372036854775807
```

(not defined for 9.223372036854776e18)

- Double comparison:

```
9223372036854775805 = 9223372036854775806 = 9223372036854775807 = 9.223372036854776e18
```

For comparison of any JSON value to SQL `NULL`, the result is `UNKNOWN`.

For comparison of JSON and non-JSON values, the non-JSON value is converted to JSON according to the rules in the following table, then the values compared as described previously.

Converting between JSON and non-JSON values

The following table provides a summary of the rules that MySQL follows when casting between JSON values and values of other types:

Table 11.2 JSON Conversion Rules

other type	CAST(other type AS JSON)	CAST(JSON AS other type)
JSON	No change	No change
utf8 character type (`utf8mb4`, `utf8`, `ascii`)	The string is parsed into a JSON value.	The JSON value is serialized into a `utf8mb4` string.
Other character types	Other character encodings are implicitly converted to `utf8mb4` and treated as described for utf8 character type.	The JSON value is serialized into a `utf8mb4` string, then cast to the other character encoding. The result may not be meaningful.
`NULL`	Results in a `NULL` value of type JSON.	Not applicable.
Geometry types	The geometry value is converted into a JSON document by calling `ST_AsGeoJSON()`.	Illegal operation. Workaround: Pass the result of `CAST(json_val AS CHAR)` to `ST_GeomFromGeoJSON()`.
All other types	Results in a JSON document consisting of a single scalar value.	Succeeds if the JSON document consists of a single scalar value of the target type and that scalar value can be cast to the target type. Otherwise, returns `NULL` and produces a warning.

`ORDER BY` and `GROUP BY` for JSON values works according to these principles:

- Ordering of scalar JSON values uses the same rules as in the preceding discussion.
- For ascending sorts, SQL `NULL` orders before all JSON values, including the JSON null literal; for descending sorts, SQL `NULL` orders after all JSON values, including the JSON null literal.
- Sort keys for JSON values are bound by the value of the `max_sort_length` system variable, so keys that differ only after the first `max_sort_length` bytes compare as equal.
- Sorting of nonscalar values is not currently supported and a warning occurs.

For sorting, it can be beneficial to cast a JSON scalar to some other native MySQL type. For example, if a column named `jdoc` contains JSON objects having a member consisting of an `id` key and a nonnegative value, use this expression to sort by `id` values:

```
ORDER BY CAST(JSON_EXTRACT(jdoc, '$.id') AS UNSIGNED)
```

If there happens to be a generated column defined to use the same expression as in the `ORDER BY`, the MySQL optimizer recognizes that and considers using the index for the query execution plan. See Section 8.3.10, “Optimizer Use of Generated Column Indexes”.

Aggregation of JSON Values

For aggregation of JSON values, SQL `NULL` values are ignored as for other data types. Non-`NULL` values are converted to a numeric type and aggregated, except for `MIN()`, `MAX()`, and `GROUP_CONCAT()`. The conversion to number should produce a meaningful result for JSON values that are numeric scalars, although (depending on the values) truncation and loss of precision may occur. Conversion to number of other JSON values may not produce a meaningful result.

11.7 Data Type Default Values

The `DEFAULT` *`value`* clause in a data type specification indicates a default value for a column. With one exception, the default value must be a constant; it cannot be a function or an expression. This means, for example, that you cannot set the default for a date column to be the value of a function such as `NOW()` or `CURRENT_DATE`. The exception is that you can specify `CURRENT_TIMESTAMP` as the default for `TIMESTAMP` and `DATETIME` columns. See Section 11.3.5, “Automatic Initialization and Updating for TIMESTAMP and DATETIME”.

`BLOB`, `TEXT`, `GEOMETRY`, and `JSON` columns cannot be assigned a default value.

If a column definition includes no explicit `DEFAULT` value, MySQL determines the default value as follows:

If the column can take `NULL` as a value, the column is defined with an explicit `DEFAULT NULL` clause.

If the column cannot take `NULL` as the value, MySQL defines the column with no explicit `DEFAULT` clause. Exception: If the column is defined as part of a `PRIMARY KEY` but not explicitly as `NOT NULL`, MySQL creates it as a `NOT NULL` column (because `PRIMARY KEY` columns must be `NOT NULL`).

For data entry into a `NOT NULL` column that has no explicit `DEFAULT` clause, if an `INSERT` or `REPLACE` statement includes no value for the column, or an `UPDATE` statement sets the column to `NULL`, MySQL handles the column according to the SQL mode in effect at the time:

- If strict SQL mode is enabled, an error occurs for transactional tables and the statement is rolled back. For nontransactional tables, an error occurs, but if this happens for the second or subsequent row of a multiple-row statement, the preceding rows will have been inserted.
- If strict mode is not enabled, MySQL sets the column to the implicit default value for the column data type.

Suppose that a table `t` is defined as follows:

```
CREATE TABLE t (i INT NOT NULL);
```

In this case, `i` has no explicit default, so in strict mode each of the following statements produce an error and no row is inserted. When not using strict mode, only the third statement produces an error; the implicit default is inserted for the first two statements, but the third fails because `DEFAULT(i)` cannot produce a value:

```
INSERT INTO t VALUES();
```

```
INSERT INTO t VALUES(DEFAULT);
INSERT INTO t VALUES(DEFAULT(i));
```

See Section 5.1.8, “Server SQL Modes”.

For a given table, you can use the `SHOW CREATE TABLE` statement to see which columns have an explicit `DEFAULT` clause.

Implicit defaults are defined as follows:

- For numeric types, the default is `0`, with the exception that for integer or floating-point types declared with the `AUTO_INCREMENT` attribute, the default is the next value in the sequence.
- For date and time types other than `TIMESTAMP`, the default is the appropriate “zero” value for the type. This is also true for `TIMESTAMP` if the `explicit_defaults_for_timestamp` system variable is enabled (see Section 5.1.5, “Server System Variables”). Otherwise, for the first `TIMESTAMP` column in a table, the default value is the current date and time. See Section 11.3, “Date and Time Types”.
- For string types other than `ENUM`, the default value is the empty string. For `ENUM`, the default is the first enumeration value.

`SERIAL DEFAULT VALUE` in the definition of an integer column is an alias for `NOT NULL AUTO_INCREMENT UNIQUE`.

11.8 Data Type Storage Requirements

- InnoDB Table Storage Requirements
- Numeric Type Storage Requirements
- Date and Time Type Storage Requirements
- String Type Storage Requirements
- Spatial Type Storage Requirements

The storage requirements for table data on disk depend on several factors. Different storage engines represent data types and store raw data differently. Table data might be compressed, either for a column or an entire row, complicating the calculation of storage requirements for a table or column.

Despite differences in storage layout on disk, the internal MySQL APIs that communicate and exchange information about table rows use a consistent data structure that applies across all storage engines.

This section includes guidelines and information for the storage requirements for each data type supported by MySQL, including the internal format and size for storage engines that use a fixed-size representation for data types. Information is listed by category or storage engine.

The internal representation of a table has a maximum row size of 65,535 bytes, even if the storage engine is capable of supporting larger rows. This figure excludes `BLOB` or `TEXT` columns, which contribute only 9 to 12 bytes toward this size. For `BLOB` and `TEXT` data, the information is stored internally in a different area of memory than the row buffer. Different storage engines handle the allocation and storage of this data in different ways, according to the method they use for handling the corresponding types. For more information, see Chapter 16, *Alternative Storage Engines*, and Section C.10.4, “Limits on Table Column Count and Row Size”.

InnoDB Table Storage Requirements

See Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table” for information about storage requirements for `InnoDB` tables.

Numeric Type Storage Requirements

Data Type	Storage Required
TINYINT	1 byte
SMALLINT	2 bytes
MEDIUMINT	3 bytes
INT, INTEGER	4 bytes
BIGINT	8 bytes
FLOAT(*p*)	4 bytes if 0 <= *p* <= 24, 8 bytes if 25 <= *p* <= 53
FLOAT	4 bytes
DOUBLE [PRECISION], REAL	8 bytes
DECIMAL(*M*,*D*), NUMERIC(*M*,*D*)	Varies; see following discussion
BIT(*M*)	approximately (*M*+7)/8 bytes

Values for DECIMAL (and NUMERIC) columns are represented using a binary format that packs nine decimal (base 10) digits into four bytes. Storage for the integer and fractional parts of each value are determined separately. Each multiple of nine digits requires four bytes, and the “leftover” digits require some fraction of four bytes. The storage required for excess digits is given by the following table.

Leftover Digits	Number of Bytes
0	0
1	1
2	1
3	2
4	2
5	3
6	3
7	4
8	4

Date and Time Type Storage Requirements

For TIME, DATETIME, and TIMESTAMP columns, the storage required for tables created before MySQL 5.6.4 differs from tables created from 5.6.4 on. This is due to a change in 5.6.4 that permits these types to have a fractional part, which requires from 0 to 3 bytes.

Data Type	Storage Required Before MySQL 5.6.4	Storage Required as of MySQL 5.6.4
YEAR	1 byte	1 byte
DATE	3 bytes	3 bytes
TIME	3 bytes	3 bytes + fractional seconds storage
DATETIME	8 bytes	5 bytes + fractional seconds storage
TIMESTAMP	4 bytes	4 bytes + fractional seconds storage

As of MySQL 5.6.4, storage for YEAR and DATE remains unchanged. However, TIME, DATETIME, and TIMESTAMP are represented differently. DATETIME is packed more efficiently, requiring 5 rather than 8 bytes for the nonfractional part, and all three parts have a fractional part that requires from 0 to 3 bytes, depending on the fractional seconds precision of stored values.

Fractional Seconds Precision	Storage Required
0	0 bytes
1, 2	1 byte
3, 4	2 bytes
5, 6	3 bytes

For example, `TIME(0)`, `TIME(2)`, `TIME(4)`, and `TIME(6)` use 3, 4, 5, and 6 bytes, respectively. `TIME` and `TIME(0)` are equivalent and require the same storage.

For details about internal representation of temporal values, see MySQL Internals: Important Algorithms and Structures.

String Type Storage Requirements

In the following table, *`M`* represents the declared column length in characters for nonbinary string types and bytes for binary string types. *`L`* represents the actual length in bytes of a given string value.

Data Type	Storage Required
`CHAR(`*`M`*`)`	*`M`* × *`w`* bytes, 0 <= *`M`* <= 255, where *`w`* is the number of bytes required for the maximum-length character in the character set. See Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table” for information about `CHAR` data type storage requirements for `InnoDB` tables.
`BINARY(`*`M`*`)`	*`M`* bytes, 0 <= *`M`* <= 255
`VARCHAR(`*`M`*`)`, `VARBINARY(`*`M`*`)`	*`L`* + 1 bytes if column values require 0 − 255 bytes, *`L`* + 2 bytes if values may require more than 255 bytes
`TINYBLOB`, `TINYTEXT`	*`L`* + 1 bytes, where *`L`* < 2^8
`BLOB`, `TEXT`	*`L`* + 2 bytes, where *`L`* < 2^{16}
`MEDIUMBLOB`, `MEDIUMTEXT`	*`L`* + 3 bytes, where *`L`* < 2^{24}
`LONGBLOB`, `LONGTEXT`	*`L`* + 4 bytes, where *`L`* < 2^{32}
`ENUM('`*`value1`*`','`*`value2`*`',...)`	1 or 2 bytes, depending on the number of enumeration values (65,535 values maximum)
`SET('`*`value1`*`','`*`value2`*`',...)`	1, 2, 3, 4, or 8 bytes, depending on the number of set members (64 members maximum)

Variable-length string types are stored using a length prefix plus data. The length prefix requires from one to four bytes depending on the data type, and the value of the prefix is *`L`* (the byte length of the string). For example, storage for a `MEDIUMTEXT` value requires *`L`* bytes to store the value plus three bytes to store the length of the value.

To calculate the number of bytes used to store a particular `CHAR`, `VARCHAR`, or `TEXT` column value, you must take into account the character set used for that column and whether the value contains multibyte characters. In particular, when using a `utf8` Unicode character set, you must keep in mind that not all characters use the same number of bytes. `utf8mb3` and `utf8mb4` character sets can require up to three and four bytes per character, respectively. For a breakdown of the storage used for different categories of `utf8mb3` or `utf8mb4` characters, see Section 10.1.9, “Unicode Support”.

`VARCHAR`, `VARBINARY`, and the `BLOB` and `TEXT` types are variable-length types. For each, the storage requirements depend on these factors:

- The actual length of the column value
- The column's maximum possible length

- The character set used for the column, because some character sets contain multibyte characters

For example, a `VARCHAR(255)` column can hold a string with a maximum length of 255 characters. Assuming that the column uses the `latin1` character set (one byte per character), the actual storage required is the length of the string (*`L`*), plus one byte to record the length of the string. For the string `'abcd'`, *`L`* is 4 and the storage requirement is five bytes. If the same column is instead declared to use the `ucs2` double-byte character set, the storage requirement is 10 bytes: The length of `'abcd'` is eight bytes and the column requires two bytes to store lengths because the maximum length is greater than 255 (up to 510 bytes).

The effective maximum number of *bytes* that can be stored in a `VARCHAR` or `VARBINARY` column is subject to the maximum row size of 65,535 bytes, which is shared among all columns. For a `VARCHAR` column that stores multibyte characters, the effective maximum number of *characters* is less. For example, `utf8mb4` characters can require up to four bytes per character, so a `VARCHAR` column that uses the `utf8mb4` character set can be declared to be a maximum of 16,383 characters. See Section C.10.4, “Limits on Table Column Count and Row Size”.

In general, the storage requirement for a `JSON` column is approximately the same as for a `LONGBLOB` or `LONGTEXT` column; that is, the space consumed by a JSON document is roughly the same as it would be for the document's string representation stored in a column of one of these types. However, there is an overhead imposed by the binary encoding, including metadata and dictionaries needed for lookup, of the individual values stored in the JSON document. For example, a string stored in a JSON document requires 4 to 10 bytes additional storage, depending on the length of the string and the size of the object or array in which it is stored.

In addition, MySQL imposes a limit on the size of any JSON document stored in a `JSON` column such that it cannot be any larger than the value of `max_allowed_packet`.

`InnoDB` encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored off-page. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.

The size of an `ENUM` object is determined by the number of different enumeration values. One byte is used for enumerations with up to 255 possible values. Two bytes are used for enumerations having between 256 and 65,535 possible values. See Section 11.4.4, “The ENUM Type”.

The size of a `SET` object is determined by the number of different set members. If the set size is *`N`*, the object occupies `(`*`N`*`+7)/8` bytes, rounded up to 1, 2, 3, 4, or 8 bytes. A `SET` can have a maximum of 64 members. See Section 11.4.5, “The SET Type”.

Spatial Type Storage Requirements

MySQL stores geometry values using 4 bytes to indicate the SRID followed by the WKB representation of the value. The `LENGTH()` function returns the space in bytes required for value storage.

For descriptions of WKB and internal storage formats for spatial values, see Section 11.5.3, “Supported Spatial Data Formats”.

11.9 Choosing the Right Type for a Column

For optimum storage, you should try to use the most precise type in all cases. For example, if an integer column is used for values in the range from `1` to `99999`, `MEDIUMINT UNSIGNED` is the best type. Of the types that represent all the required values, this type uses the least amount of storage.

All basic calculations (`+`, `-`, `*`, and `/`) with `DECIMAL` columns are done with precision of 65 decimal (base 10) digits. See Section 11.1.1, “Numeric Type Overview”.

If accuracy is not too important or if speed is the highest priority, the `DOUBLE` type may be good enough. For high precision, you can always convert to a fixed-point type stored in a `BIGINT`. This

enables you to do all calculations with 64-bit integers and then convert results back to floating-point values as necessary.

11.10 Using Data Types from Other Database Engines

To facilitate the use of code written for SQL implementations from other vendors, MySQL maps data types as shown in the following table. These mappings make it easier to import table definitions from other database systems into MySQL.

Other Vendor Type	MySQL Type
BOOL	TINYINT
BOOLEAN	TINYINT
CHARACTER VARYING(*M*)	VARCHAR(*M*)
FIXED	DECIMAL
FLOAT4	FLOAT
FLOAT8	DOUBLE
INT1	TINYINT
INT2	SMALLINT
INT3	MEDIUMINT
INT4	INT
INT8	BIGINT
LONG VARBINARY	MEDIUMBLOB
LONG VARCHAR	MEDIUMTEXT
LONG	MEDIUMTEXT
MIDDLEINT	MEDIUMINT
NUMERIC	DECIMAL

Data type mapping occurs at table creation time, after which the original type specifications are discarded. If you create a table with types used by other vendors and then issue a DESCRIBE *tbl_name* statement, MySQL reports the table structure using the equivalent MySQL types. For example:

```
mysql> CREATE TABLE t (a BOOL, b FLOAT8, c LONG VARCHAR, d NUMERIC);
Query OK, 0 rows affected (0.00 sec)

mysql> DESCRIBE t;
+-------+---------------+------+-----+---------+-------+
| Field | Type          | Null | Key | Default | Extra |
+-------+---------------+------+-----+---------+-------+
| a     | tinyint(1)    | YES  |     | NULL    |       |
| b     | double        | YES  |     | NULL    |       |
| c     | mediumtext    | YES  |     | NULL    |       |
| d     | decimal(10,0) | YES  |     | NULL    |       |
+-------+---------------+------+-----+---------+-------+
4 rows in set (0.01 sec)
```

Chapter 12 Functions and Operators

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Expressions can be used at several points in SQL statements, such as in the `ORDER BY` or `HAVING` clauses of `SELECT` statements, in the `WHERE` clause of a `SELECT`, `DELETE`, or `UPDATE` statement, or in `SET` statements. Expressions can be written using literal values, column values, `NULL`, built-in functions, stored functions, user-defined functions, and operators. This chapter describes the functions and operators that are permitted for writing expressions in MySQL. Instructions for writing stored functions and user-defined functions are given in Section 23.2, “Using Stored Routines (Procedures and Functions)”, and Section 28.4, “Adding New Functions to MySQL”. See Section 9.2.4, “Function Name Parsing and Resolution”, for the rules describing how the server interprets references to different kinds of functions.

An expression that contains `NULL` always produces a `NULL` value unless otherwise indicated in the documentation for a particular function or operator.

Note

By default, there must be no whitespace between a function name and the parenthesis following it. This helps the MySQL parser distinguish between function calls and references to tables or columns that happen to have the same name as a function. However, spaces around function arguments are permitted.

You can tell the MySQL server to accept spaces after function names by starting it with the `--sql-mode=IGNORE_SPACE` option. (See Section 5.1.8, “Server SQL Modes”.) Individual client programs can request this behavior by using the `CLIENT_IGNORE_SPACE` option for `mysql_real_connect()`. In either case, all function names become reserved words.

For the sake of brevity, most examples in this chapter display the output from the `mysql` program in abbreviated form. Rather than showing examples in this format:

```
mysql> SELECT MOD(29,9);
+-----------+
| mod(29,9) |
+-----------+
|         2 |
+-----------+
1 rows in set (0.00 sec)
```

This format is used instead:

```
mysql> SELECT MOD(29,9);
        -> 2
```

12.1 Function and Operator Reference

Table 12.1 Functions/Operators

Name	Description
ABS()	Return the absolute value
ACOS()	Return the arc cosine
ADDDATE()	Add time values (intervals) to a date value
ADDTIME()	Add time
AES_DECRYPT()	Decrypt using AES
AES_ENCRYPT()	Encrypt using AES
AND, &&	Logical AND
ANY_VALUE()	Suppress ONLY_FULL_GROUP_BY value rejection
ASCII()	Return numeric value of left-most character
ASIN()	Return the arc sine
=	Assign a value (as part of a SET statement, or as part of the SET clause in an UPDATE statement)
:=	Assign a value
ATAN()	Return the arc tangent
ATAN2(), ATAN()	Return the arc tangent of the two arguments
AVG()	Return the average value of the argument
BENCHMARK()	Repeatedly execute an expression
BETWEEN ... AND ...	Check whether a value is within a range of values
BIN()	Return a string containing binary representation of a number
BIN_TO_UUID()	Convert binary UUID to string
BINARY	Cast a string to a binary string
BIT_AND()	Return bitwise AND
BIT_COUNT()	Return the number of bits that are set
BIT_LENGTH()	Return length of argument in bits
BIT_OR()	Return bitwise OR
BIT_XOR()	Return bitwise XOR
&	Bitwise AND
~	Bitwise inversion
|	Bitwise OR
^	Bitwise XOR
CAN_ACCESS_COLUMN()	Internal use only
CAN_ACCESS_DATABASE()	Internal use only
CAN_ACCESS_TABLE()	Internal use only
CAN_ACCESS_VIEW()	Internal use only
CASE	Case operator
CAST()	Cast a value as a certain type
CEIL()	Return the smallest integer value not less than the argument
CEILING()	Return the smallest integer value not less than the argument
CHAR()	Return the character for each integer passed

Name	Description
CHAR_LENGTH()	Return number of characters in argument
CHARACTER_LENGTH()	Synonym for CHAR_LENGTH()
CHARSET()	Return the character set of the argument
COALESCE()	Return the first non-NULL argument
COERCIBILITY()	Return the collation coercibility value of the string argument
COLLATION()	Return the collation of the string argument
COMPRESS()	Return result as a binary string
CONCAT()	Return concatenated string
CONCAT_WS()	Return concatenate with separator
CONNECTION_ID()	Return the connection ID (thread ID) for the connection
CONV()	Convert numbers between different number bases
CONVERT()	Cast a value as a certain type
CONVERT_TZ()	Convert from one time zone to another
COS()	Return the cosine
COT()	Return the cotangent
COUNT()	Return a count of the number of rows returned
COUNT(DISTINCT)	Return the count of a number of different values
CRC32()	Compute a cyclic redundancy check value
CURDATE()	Return the current date
CURRENT_DATE(), CURRENT_DATE	Synonyms for CURDATE()
CURRENT_ROLE()	Returns the current active roles
CURRENT_TIME(), CURRENT_TIME	Synonyms for CURTIME()
CURRENT_TIMESTAMP(), CURRENT_TIMESTAMP	Synonyms for NOW()
CURRENT_USER(), CURRENT_USER	The authenticated user name and host name
CURTIME()	Return the current time
DATABASE()	Return the default (current) database name
DATE()	Extract the date part of a date or datetime expression
DATE_ADD()	Add time values (intervals) to a date value
DATE_FORMAT()	Format date as specified
DATE_SUB()	Subtract a time value (interval) from a date
DATEDIFF()	Subtract two dates
DAY()	Synonym for DAYOFMONTH()
DAYNAME()	Return the name of the weekday
DAYOFMONTH()	Return the day of the month (0-31)
DAYOFWEEK()	Return the weekday index of the argument
DAYOFYEAR()	Return the day of the year (1-366)
DECODE()	Decodes a string encrypted using ENCODE()
DEFAULT()	Return the default value for a table column
DEGREES()	Convert radians to degrees
DES_DECRYPT()	Decrypt a string

Name	Description
DES_ENCRYPT()	Encrypt a string
DIV	Integer division
/	Division operator
ELT()	Return string at index number
ENCODE()	Encode a string
ENCRYPT()	Encrypt a string
=	Equal operator
<=>	NULL-safe equal to operator
EXP()	Raise to the power of
EXPORT_SET()	Return a string such that for every bit set in the value bits, you get an on string and for every unset bit, you get an off string
EXTRACT()	Extract part of a date
ExtractValue()	Extracts a value from an XML string using XPath notation
FIELD()	Return the index (position) of the first argument in the subsequent arguments
FIND_IN_SET()	Return the index position of the first argument within the second argument
FLOOR()	Return the largest integer value not greater than the argument
FORMAT()	Return a number formatted to specified number of decimal places
FOUND_ROWS()	For a SELECT with a LIMIT clause, the number of rows that would be returned were there no LIMIT clause
FROM_BASE64()	Decode to a base-64 string and return result
FROM_DAYS()	Convert a day number to a date
FROM_UNIXTIME()	Format Unix timestamp as a date
GeometryCollection()	Construct geometry collection from geometries
GET_DD_COLUMN_PRIVILEGES()	Internal use only
GET_DD_CREATE_OPTIONS()	Internal use only
GET_DD_INDEX_SUB_PART_LENGTH()	Internal use only
GET_FORMAT()	Return a date format string
GET_LOCK()	Get a named lock
>	Greater than operator
>=	Greater than or equal operator
GREATEST()	Return the largest argument
GROUP_CONCAT()	Return a concatenated string
GROUPING()	Distinguish super-aggregate ROLLUP rows from regular rows
GTID_SUBSET()	Return true if all GTIDs in subset are also in set; otherwise false.
GTID_SUBTRACT()	Return all GTIDs in set that are not in subset.
HEX()	Return a hexadecimal representation of a decimal or string value

Name	Description
HOUR()	Extract the hour
IF()	If/else construct
IFNULL()	Null if/else construct
IN()	Check whether a value is within a set of values
INET_ATON()	Return the numeric value of an IP address
INET_NTOA()	Return the IP address from a numeric value
INET6_ATON()	Return the numeric value of an IPv6 address
INET6_NTOA()	Return the IPv6 address from a numeric value
INSERT()	Insert a substring at the specified position up to the specified number of characters
INSTR()	Return the index of the first occurrence of substring
INTERNAL_AUTO_INCREMENT()	Internal use only
INTERNAL_AVG_ROW_LENGTH()	Internal use only
INTERNAL_CHECK_TIME()	Internal use only
INTERNAL_CHECKSUM()	Internal use only
INTERNAL_DATA_FREE()	Internal use only
INTERNAL_DATA_LENGTH()	Internal use only
INTERNAL_DD_CHAR_LENGTH()	Internal use only
INTERNAL_GET_COMMENT_OR_ERROR()	Internal use only
INTERNAL_GET_VIEW_WARNING_OR_ERROR()	Internal use only
INTERNAL_INDEX_COLUMN_CARDINALITY()	Internal use only
INTERNAL_INDEX_LENGTH()	Internal use only
INTERNAL_KEYS_DISABLED()	Internal use only
INTERNAL_MAX_DATA_LENGTH()	Internal use only
INTERNAL_TABLE_ROWS()	Internal use only
INTERNAL_UPDATE_TIME()	Internal use only
INTERVAL()	Return the index of the argument that is less than the first argument
IS	Test a value against a boolean
IS_FREE_LOCK()	Whether the named lock is free
IS_IPV4()	Whether argument is an IPv4 address
IS_IPV4_COMPAT()	Whether argument is an IPv4-compatible address
IS_IPV4_MAPPED()	Whether argument is an IPv4-mapped address
IS_IPV6()	Whether argument is an IPv6 address
IS NOT	Test a value against a boolean
IS NOT NULL	NOT NULL value test
IS NULL	NULL value test
IS_USED_LOCK()	Whether the named lock is in use; return connection identifier if true
IS_UUID()	Whether argument is a valid UUID
ISNULL()	Test whether the argument is NULL
JSON_ARRAY()	Create JSON array

Name	Description
JSON_ARRAY_APPEND()	Append data to JSON document
JSON_ARRAY_INSERT()	Insert into JSON array
JSON_ARRAYAGG()	Return result set as a single JSON array
->	Return value from JSON column after evaluating path; equivalent to JSON_EXTRACT().
JSON_CONTAINS()	Whether JSON document contains specific object at path
JSON_CONTAINS_PATH()	Whether JSON document contains any data at path
JSON_DEPTH()	Maximum depth of JSON document
JSON_EXTRACT()	Return data from JSON document
->>	Return value from JSON column after evaluating path and unquoting the result; equivalent to JSON_UNQUOTE(JSON_EXTRACT()).
JSON_INSERT()	Insert data into JSON document
JSON_KEYS()	Array of keys from JSON document
JSON_LENGTH()	Number of elements in JSON document
JSON_MERGE()	Merge JSON documents
JSON_OBJECT()	Create JSON object
JSON_OBJECTAGG()	Return result set as a single JSON object
JSON_PRETTY()	Prints a JSON document in human-readable format, with each array element or object member printed on a new line, indented two spaces with respect to its parent.
JSON_QUOTE()	Quote JSON document
JSON_REMOVE()	Remove data from JSON document
JSON_REPLACE()	Replace values in JSON document
JSON_SEARCH()	Path to value within JSON document
JSON_SET()	Insert data into JSON document
JSON_STORAGE_FREE()	Freed space within binary representation of a JSON column value following a partial update
JSON_STORAGE_SIZE()	Space used for storage of binary representation of a JSON document; for a JSON column, the space used when the document was inserted, prior to any partial updates
JSON_TYPE()	Type of JSON value
JSON_UNQUOTE()	Unquote JSON value
JSON_VALID()	Whether JSON value is valid
LAST_DAY	Return the last day of the month for the argument
LAST_INSERT_ID()	Value of the AUTOINCREMENT column for the last INSERT
LCASE()	Synonym for LOWER()
LEAST()	Return the smallest argument
LEFT()	Return the leftmost number of characters as specified
<<	Left shift
LENGTH()	Return the length of a string in bytes
<	Less than operator
<=	Less than or equal operator

Name	Description
LIKE	Simple pattern matching
LineString()	Construct LineString from Point values
LN()	Return the natural logarithm of the argument
LOAD_FILE()	Load the named file
LOCALTIME(), LOCALTIME	Synonym for NOW()
LOCALTIMESTAMP, LOCALTIMESTAMP()	Synonym for NOW()
LOCATE()	Return the position of the first occurrence of substring
LOG()	Return the natural logarithm of the first argument
LOG10()	Return the base-10 logarithm of the argument
LOG2()	Return the base-2 logarithm of the argument
LOWER()	Return the argument in lowercase
LPAD()	Return the string argument, left-padded with the specified string
LTRIM()	Remove leading spaces
MAKE_SET()	Return a set of comma-separated strings that have the corresponding bit in bits set
MAKEDATE()	Create a date from the year and day of year
MAKETIME()	Create time from hour, minute, second
MASTER_POS_WAIT()	Block until the slave has read and applied all updates up to the specified position
MATCH	Perform full-text search
MAX()	Return the maximum value
MBRContains()	Whether MBR of one geometry contains MBR of another
MBRCoveredBy()	Whether one MBR is covered by another
MBRCovers()	Whether one MBR covers another
MBRDisjoint()	Whether MBRs of two geometries are disjoint
MBREquals()	Whether MBRs of two geometries are equal
MBRIntersects()	Whether MBRs of two geometries intersect
MBROverlaps()	Whether MBRs of two geometries overlap
MBRTouches()	Whether MBRs of two geometries touch
MBRWithin()	Whether MBR of one geometry is within MBR of another
MD5()	Calculate MD5 checksum
MICROSECOND()	Return the microseconds from argument
MID()	Return a substring starting from the specified position
MIN()	Return the minimum value
-	Minus operator
MINUTE()	Return the minute from the argument
MOD()	Return the remainder
%, MOD	Modulo operator
MONTH()	Return the month from the date passed
MONTHNAME()	Return the name of the month

Name	Description
MultiLineString()	Contruct MultiLineString from LineString values
MultiPoint()	Construct MultiPoint from Point values
MultiPolygon()	Construct MultiPolygon from Polygon values
NAME_CONST()	Causes the column to have the given name
NOT, !	Negates value
NOT BETWEEN ... AND ...	Check whether a value is not within a range of values
!=, <>	Not equal operator
NOT IN()	Check whether a value is not within a set of values
NOT LIKE	Negation of simple pattern matching
NOT REGEXP	Negation of REGEXP
NOW()	Return the current date and time
NULLIF()	Return NULL if expr1 = expr2
OCT()	Return a string containing octal representation of a number
OCTET_LENGTH()	Synonym for LENGTH()
||, OR	Logical OR
ORD()	Return character code for leftmost character of the argument
PASSWORD()	Calculate and return a password string
PERIOD_ADD()	Add a period to a year-month
PERIOD_DIFF()	Return the number of months between periods
PI()	Return the value of pi
+	Addition operator
Point()	Construct Point from coordinates
Polygon()	Construct Polygon from LineString arguments
POSITION()	Synonym for LOCATE()
POW()	Return the argument raised to the specified power
POWER()	Return the argument raised to the specified power
QUARTER()	Return the quarter from a date argument
QUOTE()	Escape the argument for use in an SQL statement
RADIANS()	Return argument converted to radians
RAND()	Return a random floating-point value
RANDOM_BYTES()	Return a random byte vector
REGEXP	Pattern matching using regular expressions
RELEASE_ALL_LOCKS()	Releases all current named locks
RELEASE_LOCK()	Releases the named lock
REPEAT()	Repeat a string the specified number of times
REPLACE()	Replace occurrences of a specified string
REVERSE()	Reverse the characters in a string
RIGHT()	Return the specified rightmost number of characters
>>	Right shift
RLIKE	Synonym for REGEXP

Name	Description
ROLES_GRAPHML()	Returns a GraphML document representing memory role subgraphs
ROUND()	Round the argument
ROW_COUNT()	The number of rows updated
RPAD()	Append string the specified number of times
RTRIM()	Remove trailing spaces
SCHEMA()	Synonym for DATABASE()
SEC_TO_TIME()	Converts seconds to 'HH:MM:SS' format
SECOND()	Return the second (0-59)
SESSION_USER()	Synonym for USER()
SHA1(), SHA()	Calculate an SHA-1 160-bit checksum
SHA2()	Calculate an SHA-2 checksum
SIGN()	Return the sign of the argument
SIN()	Return the sine of the argument
SLEEP()	Sleep for a number of seconds
SOUNDEX()	Return a soundex string
SOUNDS LIKE	Compare sounds
SPACE()	Return a string of the specified number of spaces
SQRT()	Return the square root of the argument
ST_Area()	Return Polygon or MultiPolygon area
ST_AsBinary(), ST_AsWKB()	Convert from internal geometry format to WKB
ST_AsGeoJSON()	Generate GeoJSON object from geometry
ST_AsText(), ST_AsWKT()	Convert from internal geometry format to WKT
ST_Buffer()	Return geometry of points within given distance from geometry
ST_Buffer_Strategy()	Produce strategy option for ST_Buffer()
ST_Centroid()	Return centroid as a point
ST_Contains()	Whether one geometry contains another
ST_ConvexHull()	Return convex hull of geometry
ST_Crosses()	Whether one geometry crosses another
ST_Difference()	Return point set difference of two geometries
ST_Dimension()	Dimension of geometry
ST_Disjoint()	Whether one geometry is disjoint from another
ST_Distance()	The distance of one geometry from another
ST_Distance_Sphere()	Minimum distance on earth between two geometries
ST_EndPoint()	End Point of LineString
ST_Envelope()	Return MBR of geometry
ST_Equals()	Whether one geometry is equal to another
ST_ExteriorRing()	Return exterior ring of Polygon
ST_GeoHash()	Produce a geohash value

Name	Description
ST_GeomCollFromText(), ST_GeometryCollectionFromText(), ST_GeomCollFromTxt()	Return geometry collection from WKT
ST_GeomCollFromWKB(), ST_GeometryCollectionFromWKB()	Return geometry collection from WKB
ST_GeometryN()	Return N-th geometry from geometry collection
ST_GeometryType()	Return name of geometry type
ST_GeomFromGeoJSON()	Generate geometry from GeoJSON object
ST_GeomFromText(), ST_GeometryFromText()	Return geometry from WKT
ST_GeomFromWKB(), ST_GeometryFromWKB()	Return geometry from WKB
ST_InteriorRingN()	Return N-th interior ring of Polygon
ST_Intersection()	Return point set intersection of two geometries
ST_Intersects()	Whether one geometry intersects another
ST_IsClosed()	Whether a geometry is closed and simple
ST_IsEmpty()	Placeholder function
ST_IsSimple()	Whether a geometry is simple
ST_IsValid()	Whether a geometry is valid
ST_LatFromGeoHash()	Return latitude from geohash value
ST_Length()	Return length of LineString
ST_LineFromText(), ST_LineStringFromText()	Construct LineString from WKT
ST_LineFromWKB(), ST_LineStringFromWKB()	Construct LineString from WKB
ST_LongFromGeoHash()	Return longitude from geohash value
ST_MakeEnvelope()	Rectangle around two points
ST_MLineFromText(), ST_MultiLineStringFromText()	Construct MultiLineString from WKT
ST_MLineFromWKB(), ST_MultiLineStringFromWKB()	Construct MultiLineString from WKB
ST_MPointFromText(), ST_MultiPointFromText()	Construct MultiPoint from WKT
ST_MPointFromWKB(), ST_MultiPointFromWKB()	Construct MultiPoint from WKB
ST_MPolyFromText(), ST_MultiPolygonFromText()	Construct MultiPolygon from WKT
ST_MPolyFromWKB(), ST_MultiPolygonFromWKB()	Construct MultiPolygon from WKB
ST_NumGeometries()	Return number of geometries in geometry collection
ST_NumInteriorRing(), ST_NumInteriorRings()	Return number of interior rings in Polygon
ST_NumPoints()	Return number of points in LineString
ST_Overlaps()	Whether one geometry overlaps another
ST_PointFromGeoHash()	Convert geohash value to POINT value

Name	Description
ST_PointFromText()	Construct Point from WKT
ST_PointFromWKB()	Construct Point from WKB
ST_PointN()	Return N-th point from LineString
ST_PolyFromText(), ST_PolygonFromText()	Construct Polygon from WKT
ST_PolyFromWKB(), ST_PolygonFromWKB()	Construct Polygon from WKB
ST_Simplify()	Return simplified geometry
ST_SRID()	Return spatial reference system ID for geometry
ST_StartPoint()	Start Point of LineString
ST_SwapXY()	Return argument with X/Y coordinates swapped
ST_SymDifference()	Return point set symmetric difference of two geometries
ST_Touches()	Whether one geometry touches another
ST_Union()	Return point set union of two geometries
ST_Validate()	Return validated geometry
ST_Within()	Whether one geometry is within another
ST_X()	Return X coordinate of Point
ST_Y()	Return Y coordinate of Point
STD()	Return the population standard deviation
STDDEV()	Return the population standard deviation
STDDEV_POP()	Return the population standard deviation
STDDEV_SAMP()	Return the sample standard deviation
STR_TO_DATE()	Convert a string to a date
STRCMP()	Compare two strings
SUBDATE()	Synonym for DATE_SUB() when invoked with three arguments
SUBSTR()	Return the substring as specified
SUBSTRING()	Return the substring as specified
SUBSTRING_INDEX()	Return a substring from a string before the specified number of occurrences of the delimiter
SUBTIME()	Subtract times
SUM()	Return the sum
SYSDATE()	Return the time at which the function executes
SYSTEM_USER()	Synonym for USER()
TAN()	Return the tangent of the argument
TIME()	Extract the time portion of the expression passed
TIME_FORMAT()	Format as time
TIME_TO_SEC()	Return the argument converted to seconds
TIMEDIFF()	Subtract time
*	Multiplication operator

Name	Description
TIMESTAMP()	With a single argument, this function returns the date or datetime expression; with two arguments, the sum of the arguments
TIMESTAMPADD()	Add an interval to a datetime expression
TIMESTAMPDIFF()	Subtract an interval from a datetime expression
TO_BASE64()	Return the argument converted to a base-64 string
TO_DAYS()	Return the date argument converted to days
TO_SECONDS()	Return the date or datetime argument converted to seconds since Year 0
TRIM()	Remove leading and trailing spaces
TRUNCATE()	Truncate to specified number of decimal places
UCASE()	Synonym for UPPER()
-	Change the sign of the argument
UNCOMPRESS()	Uncompress a string compressed
UNCOMPRESSED_LENGTH()	Return the length of a string before compression
UNHEX()	Return a string containing hex representation of a number
UNIX_TIMESTAMP()	Return a Unix timestamp
UpdateXML()	Return replaced XML fragment
UPPER()	Convert to uppercase
USER()	The user name and host name provided by the client
UTC_DATE()	Return the current UTC date
UTC_TIME()	Return the current UTC time
UTC_TIMESTAMP()	Return the current UTC date and time
UUID()	Return a Universal Unique Identifier (UUID)
UUID_SHORT()	Return an integer-valued universal identifier
UUID_TO_BIN()	Convert string UUID to binary
VALIDATE_PASSWORD_STRENGTH()	Determine strength of password
VALUES()	Defines the values to be used during an INSERT
VAR_POP()	Return the population standard variance
VAR_SAMP()	Return the sample variance
VARIANCE()	Return the population standard variance
VERSION()	Return a string that indicates the MySQL server version
WAIT_FOR_EXECUTED_GTID_SET()	Wait until the given GTIDs have executed on slave.
WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()	Wait until the given GTIDs have executed on slave.
WEEK()	Return the week number
WEEKDAY()	Return the weekday index
WEEKOFYEAR()	Return the calendar week of the date (1-53)
WEIGHT_STRING()	Return the weight string for a string
XOR	Logical XOR
YEAR()	Return the year
YEARWEEK()	Return the year and week

12.2 Type Conversion in Expression Evaluation

When an operator is used with operands of different types, type conversion occurs to make the operands compatible. Some conversions occur implicitly. For example, MySQL automatically converts numbers to strings as necessary, and vice versa.

```
mysql> SELECT 1+'1';
        -> 2
mysql> SELECT CONCAT(2,' test');
        -> '2 test'
```

It is also possible to convert a number to a string explicitly using the `CAST()` function. Conversion occurs implicitly with the `CONCAT()` function because it expects string arguments.

```
mysql> SELECT 38.8, CAST(38.8 AS CHAR);
        -> 38.8, '38.8'
mysql> SELECT 38.8, CONCAT(38.8);
        -> 38.8, '38.8'
```

See later in this section for information about the character set of implicit number-to-string conversions, and for modified rules that apply to `CREATE TABLE ... SELECT` statements.

The following rules describe how conversion occurs for comparison operations:

- If one or both arguments are `NULL`, the result of the comparison is `NULL`, except for the `NULL`-safe `<=>` equality comparison operator. For `NULL <=> NULL`, the result is true. No conversion is needed.
- If both arguments in a comparison operation are strings, they are compared as strings.
- If both arguments are integers, they are compared as integers.
- Hexadecimal values are treated as binary strings if not compared to a number.
- If one of the arguments is a `TIMESTAMP` or `DATETIME` column and the other argument is a constant, the constant is converted to a timestamp before the comparison is performed. This is done to be more ODBC-friendly. Note that this is not done for the arguments to `IN()`! To be safe, always use complete datetime, date, or time strings when doing comparisons. For example, to achieve best results when using `BETWEEN` with date or time values, use `CAST()` to explicitly convert the values to the desired data type.

  A single-row subquery from a table or tables is not considered a constant. For example, if a subquery returns an integer to be compared to a `DATETIME` value, the comparison is done as two integers. The integer is not converted to a temporal value. To compare the operands as `DATETIME` values, use `CAST()` to explicitly convert the subquery value to `DATETIME`.
- If one of the arguments is a decimal value, comparison depends on the other argument. The arguments are compared as decimal values if the other argument is a decimal or integer value, or as floating-point values if the other argument is a floating-point value.
- In all other cases, the arguments are compared as floating-point (real) numbers.

For information about conversion of values from one temporal type to another, see Section 11.3.7, “Conversion Between Date and Time Types”.

Comparison of JSON values takes place at two levels. The first level of comparison is based on the JSON types of the compared values. If the types differ, the comparison result is determined solely by which type has higher precedence. If the two values have the same JSON type, a second level of comparison occurs using type-specific rules. For comparison of JSON and non-JSON values, the non-JSON value is converted to JSON and the values compared as JSON values. For details, see Comparison and Ordering of JSON Values.

The following examples illustrate conversion of strings to numbers for comparison operations:

```
mysql> SELECT 1 > '6x';
        -> 0
mysql> SELECT 7 > '6x';
        -> 1
mysql> SELECT 0 > 'x6';
        -> 0
mysql> SELECT 0 = 'x6';
        -> 1
```

For comparisons of a string column with a number, MySQL cannot use an index on the column to look up the value quickly. If *`str_col`* is an indexed string column, the index cannot be used when performing the lookup in the following statement:

```
SELECT * FROM tbl_name WHERE str_col=1;
```

The reason for this is that there are many different strings that may convert to the value `1`, such as `'1'`, `' 1'`, or `'1a'`.

Comparisons that use floating-point numbers (or values that are converted to floating-point numbers) are approximate because such numbers are inexact. This might lead to results that appear inconsistent:

```
mysql> SELECT '18015376320243458' = 18015376320243458;
        -> 1
mysql> SELECT '18015376320243459' = 18015376320243459;
        -> 0
```

Such results can occur because the values are converted to floating-point numbers, which have only 53 bits of precision and are subject to rounding:

```
mysql> SELECT '18015376320243459'+0.0;
        -> 1.8015376320243e+16
```

Furthermore, the conversion from string to floating-point and from integer to floating-point do not necessarily occur the same way. The integer may be converted to floating-point by the CPU, whereas the string is converted digit by digit in an operation that involves floating-point multiplications.

The results shown will vary on different systems, and can be affected by factors such as computer architecture or the compiler version or optimization level. One way to avoid such problems is to use `CAST()` so that a value is not converted implicitly to a float-point number:

```
mysql> SELECT CAST('18015376320243459' AS UNSIGNED) = 18015376320243459;
        -> 1
```

For more information about floating-point comparisons, see Section B.5.4.8, “Problems with Floating-Point Values”.

The server includes `dtoa`, a conversion library that provides the basis for improved conversion between string or `DECIMAL` values and approximate-value (`FLOAT`/`DOUBLE`) numbers:

- Consistent conversion results across platforms, which eliminates, for example, Unix versus Windows conversion differences.
- Accurate representation of values in cases where results previously did not provide sufficient precision, such as for values close to IEEE limits.
- Conversion of numbers to string format with the best possible precision. The precision of `dtoa` is always the same or better than that of the standard C library functions.

Because the conversions produced by this library differ in some cases from non-`dtoa` results, the potential exists for incompatibilities in applications that rely on previous results. For example,

applications that depend on a specific exact result from previous conversions might need adjustment to accommodate additional precision.

The `dtoa` library provides conversions with the following properties. *`D`* represents a value with a `DECIMAL` or string representation, and *`F`* represents a floating-point number in native binary (IEEE) format.

- *`F`* -> *`D`* conversion is done with the best possible precision, returning *`D`* as the shortest string that yields *`F`* when read back in and rounded to the nearest value in native binary format as specified by IEEE.
- *`D`* -> *`F`* conversion is done such that *`F`* is the nearest native binary number to the input decimal string *`D`*.

These properties imply that *`F`* -> *`D`* -> *`F`* conversions are lossless unless *`F`* is `-inf`, `+inf`, or `NaN`. The latter values are not supported because the SQL standard defines them as invalid values for `FLOAT` or `DOUBLE`.

For *`D`* -> *`F`* -> *`D`* conversions, a sufficient condition for losslessness is that *`D`* uses 15 or fewer digits of precision, is not a denormal value, `-inf`, `+inf`, or `NaN`. In some cases, the conversion is lossless even if *`D`* has more than 15 digits of precision, but this is not always the case.

Implicit conversion of a numeric or temporal value to string produces a value that has a character set and collation determined by the `character_set_connection` and `collation_connection` system variables. (These variables commonly are set with `SET NAMES`. For information about connection character sets, see Section 10.1.4, “Connection Character Sets and Collations”.)

This means that such a conversion results in a character (nonbinary) string (a `CHAR`, `VARCHAR`, or `LONGTEXT` value), except in the case that the connection character set is set to `binary`. In that case, the conversion result is a binary string (a `BINARY`, `VARBINARY`, or `LONGBLOB` value).

For integer expressions, the preceding remarks about expression *evaluation* apply somewhat differently for expression *assignment*; for example, in a statement such as this:

```
CREATE TABLE t SELECT integer_expr;
```

In this case, the table in the column resulting from the expression has type `INT` or `BIGINT` depending on the length of the integer expression. If the maximum length of the expression does not fit in an `INT`, `BIGINT` is used instead. The length is taken from the `max_length` value of the `SELECT` result set metadata (see Section 27.7.5, “C API Data Structures”). This means that you can force a `BIGINT` rather than `INT` by use of a sufficiently long expression:

```
CREATE TABLE t SELECT 00000000000000000000;
```

12.3 Operators

Table 12.2 Operators

Name	Description
`AND`, `&&`	Logical AND
`=`	Assign a value (as part of a `SET` statement, or as part of the `SET` clause in an `UPDATE` statement)
`:=`	Assign a value
`BETWEEN ... AND ...`	Check whether a value is within a range of values
`BINARY`	Cast a string to a binary string
`&`	Bitwise AND
`~`	Bitwise inversion
`|`	Bitwise OR

Name	Description
^	Bitwise XOR
CASE	Case operator
DIV	Integer division
/	Division operator
=	Equal operator
<=>	NULL-safe equal to operator
>	Greater than operator
>=	Greater than or equal operator
IS	Test a value against a boolean
IS NOT	Test a value against a boolean
IS NOT NULL	NOT NULL value test
IS NULL	NULL value test
->	Return value from JSON column after evaluating path; equivalent to JSON_EXTRACT().
->>	Return value from JSON column after evaluating path and unquoting the result; equivalent to JSON_UNQUOTE(JSON_EXTRACT()).
<<	Left shift
<	Less than operator
<=	Less than or equal operator
LIKE	Simple pattern matching
-	Minus operator
%, MOD	Modulo operator
NOT, !	Negates value
NOT BETWEEN ... AND ...	Check whether a value is not within a range of values
!=, <>	Not equal operator
NOT LIKE	Negation of simple pattern matching
NOT REGEXP	Negation of REGEXP
||, OR	Logical OR
+	Addition operator
REGEXP	Pattern matching using regular expressions
>>	Right shift
RLIKE	Synonym for REGEXP
SOUNDS LIKE	Compare sounds
*	Multiplication operator
-	Change the sign of the argument
XOR	Logical XOR

12.3.1 Operator Precedence

Operator precedences are shown in the following list, from highest precedence to the lowest. Operators that are shown together on a line have the same precedence.

```
INTERVAL
```

```
BINARY, COLLATE
!
- (unary minus), ~ (unary bit inversion)
^
*, /, DIV, %, MOD
-, +
<<, >>
&
|
= (comparison), <=>, >=, >, <=, <, <>, !=, IS, LIKE, REGEXP, IN
BETWEEN, CASE, WHEN, THEN, ELSE
NOT
AND, &&
XOR
OR, ||
= (assignment), :=
```

The precedence of `=` depends on whether it is used as a comparison operator (`=`) or as an assignment operator (`=`). When used as a comparison operator, it has the same precedence as `<=>`, `>=`, `>`, `<=`, `<`, `<>`, `!=`, `IS`, `LIKE`, `REGEXP`, and `IN`. When used as an assignment operator, it has the same precedence as `:=`. Section 13.7.4.1, “SET Syntax for Variable Assignment”, and Section 9.4, “User-Defined Variables”, explain how MySQL determines which interpretation of `=` should apply.

For operators that occur at the same precedence level within an expression, evaluation proceeds left to right, with the exception that assignments evaluate right to left.

The meaning of some operators depends on the SQL mode:

- By default, `||` is a logical `OR` operator. With `PIPES_AS_CONCAT` enabled, `||` is string concatenation, with a precedence between `^` and the unary operators.
- By default, `!` has a higher precedence than `NOT`. With `HIGH_NOT_PRECEDENCE` enabled, `!` and `NOT` have the same precedence.

See Section 5.1.8, “Server SQL Modes”.

The precedence of operators determines the order of evaluation of terms in an expression. To override this order and group terms explicitly, use parentheses. For example:

```
mysql> SELECT 1+2*3;
        -> 7
mysql> SELECT (1+2)*3;
        -> 9
```

12.3.2 Comparison Functions and Operators

Table 12.3 Comparison Operators

Name	Description
`BETWEEN ... AND ...`	Check whether a value is within a range of values
`COALESCE()`	Return the first non-NULL argument
`=`	Equal operator
`<=>`	NULL-safe equal to operator
`>`	Greater than operator
`>=`	Greater than or equal operator
`GREATEST()`	Return the largest argument
`IN()`	Check whether a value is within a set of values
`INTERVAL()`	Return the index of the argument that is less than the first argument
`IS`	Test a value against a boolean

Name	Description
`IS NOT`	Test a value against a boolean
`IS NOT NULL`	NOT NULL value test
`IS NULL`	NULL value test
`ISNULL()`	Test whether the argument is NULL
`LEAST()`	Return the smallest argument
`<`	Less than operator
`<=`	Less than or equal operator
`LIKE`	Simple pattern matching
`NOT BETWEEN ... AND ...`	Check whether a value is not within a range of values
`!=, <>`	Not equal operator
`NOT IN()`	Check whether a value is not within a set of values
`NOT LIKE`	Negation of simple pattern matching
`STRCMP()`	Compare two strings

Comparison operations result in a value of `1` (`TRUE`), `0` (`FALSE`), or `NULL`. These operations work for both numbers and strings. Strings are automatically converted to numbers and numbers to strings as necessary.

The following relational comparison operators can be used to compare not only scalar operands, but row operands:

```
=  >  <  >=  <=  <>  !=
```

The descriptions for those operators later in this section detail how they work with row operands. For additional examples of row comparisons in the context of row subqueries, see Section 13.2.11.5, “Row Subqueries”.

Some of the functions in this section return values other than `1` (`TRUE`), `0` (`FALSE`), or `NULL`. For example, `LEAST()` and `GREATEST()`. However, the value they return is based on comparison operations performed according to the rules described in Section 12.2, “Type Conversion in Expression Evaluation”.

To convert a value to a specific type for comparison purposes, you can use the `CAST()` function. String values can be converted to a different character set using `CONVERT()`. See Section 12.10, “Cast Functions and Operators”.

By default, string comparisons are not case sensitive and use the current character set. The default is `utf8mb4`.

- `=`

  Equal:

```
mysql> SELECT 1 = 0;
        -> 0
mysql> SELECT '0' = 0;
        -> 1
mysql> SELECT '0.0' = 0;
        -> 1
mysql> SELECT '0.01' = 0;
        -> 0
mysql> SELECT '.01' = 0.01;
        -> 1
```

  For row comparisons, `(a, b) = (x, y)` is equivalent to:

```
(a = x) AND (b = y)
```

- `<=>`

  `NULL`-safe equal. This operator performs an equality comparison like the `=` operator, but returns `1` rather than `NULL` if both operands are `NULL`, and `0` rather than `NULL` if one operand is `NULL`.

  The `<=>` operator is equivalent to the standard SQL `IS NOT DISTINCT FROM` operator.

```
mysql> SELECT 1 <=> 1, NULL <=> NULL, 1 <=> NULL;
        -> 1, 1, 0
mysql> SELECT 1 = 1, NULL = NULL, 1 = NULL;
        -> 1, NULL, NULL
```

  For row comparisons, `(a, b) <=> (x, y)` is equivalent to:

```
(a <=> x) AND (b <=> y)
```

- `<>`, `!=`

  Not equal:

```
mysql> SELECT '.01' <> '0.01';
        -> 1
mysql> SELECT .01 <> '0.01';
        -> 0
mysql> SELECT 'zapp' <> 'zappp';
        -> 1
```

  For row comparisons, `(a, b) <> (x, y)` and `(a, b) != (x, y)` are equivalent to:

```
(a <> x) OR (b <> y)
```

- `<=`

  Less than or equal:

```
mysql> SELECT 0.1 <= 2;
        -> 1
```

  For row comparisons, `(a, b) <= (x, y)` is equivalent to:

```
(a < x) OR ((a = x) AND (b <= y))
```

- `<`

  Less than:

```
mysql> SELECT 2 < 2;
        -> 0
```

  For row comparisons, `(a, b) < (x, y)` is equivalent to:

```
(a < x) OR ((a = x) AND (b < y))
```

- `>=`

  Greater than or equal:

```
mysql> SELECT 2 >= 2;
       -> 1
```

For row comparisons, `(a, b) >= (x, y)` is equivalent to:

```
(a > x) OR ((a = x) AND (b >= y))
```

- `>`

  Greater than:

  ```
  mysql> SELECT 2 > 2;
         -> 0
  ```

  For row comparisons, `(a, b) > (x, y)` is equivalent to:

  ```
  (a > x) OR ((a = x) AND (b > y))
  ```

- `IS boolean_value`

  Tests a value against a boolean value, where *`boolean_value`* can be `TRUE`, `FALSE`, or `UNKNOWN`.

  ```
  mysql> SELECT 1 IS TRUE, 0 IS FALSE, NULL IS UNKNOWN;
         -> 1, 1, 1
  ```

- `IS NOT boolean_value`

  Tests a value against a boolean value, where *`boolean_value`* can be `TRUE`, `FALSE`, or `UNKNOWN`.

  ```
  mysql> SELECT 1 IS NOT UNKNOWN, 0 IS NOT UNKNOWN, NULL IS NOT UNKNOWN;
         -> 1, 1, 0
  ```

- `IS NULL`

  Tests whether a value is `NULL`.

  ```
  mysql> SELECT 1 IS NULL, 0 IS NULL, NULL IS NULL;
         -> 0, 0, 1
  ```

  To work well with ODBC programs, MySQL supports the following extra features when using `IS NULL`:

  - If `sql_auto_is_null` variable is set to 1, then after a statement that successfully inserts an automatically generated `AUTO_INCREMENT` value, you can find that value by issuing a statement of the following form:

    ```
    SELECT * FROM tbl_name WHERE auto_col IS NULL
    ```

    If the statement returns a row, the value returned is the same as if you invoked the `LAST_INSERT_ID()` function. For details, including the return value after a multiple-row insert, see Section 12.14, “Information Functions”. If no `AUTO_INCREMENT` value was successfully inserted, the `SELECT` statement returns no row.

    The behavior of retrieving an `AUTO_INCREMENT` value by using an `IS NULL` comparison can be disabled by setting `sql_auto_is_null = 0`. See Section 5.1.5, “Server System Variables”.

    The default value of `sql_auto_is_null` is 0.

- For `DATE` and `DATETIME` columns that are declared as `NOT NULL`, you can find the special date `'0000-00-00'` by using a statement like this:

```
SELECT * FROM tbl_name WHERE date_column IS NULL
```

  This is needed to get some ODBC applications to work because ODBC does not support a `'0000-00-00'` date value.

  See Obtaining Auto-Increment Values, and the description for the `FLAG_AUTO_IS_NULL` option at Connector/ODBC Connection Parameters.

- `IS NOT NULL`

  Tests whether a value is not `NULL`.

```
mysql> SELECT 1 IS NOT NULL, 0 IS NOT NULL, NULL IS NOT NULL;
        -> 1, 1, 0
```

- *`expr`* `BETWEEN` *`min`* `AND` *`max`*

  If *`expr`* is greater than or equal to *`min`* and *`expr`* is less than or equal to *`max`*, `BETWEEN` returns `1`, otherwise it returns `0`. This is equivalent to the expression `(`*`min`* `<=` *`expr`* `AND` *`expr`* `<=` *`max`*`)` if all the arguments are of the same type. Otherwise type conversion takes place according to the rules described in Section 12.2, “Type Conversion in Expression Evaluation”, but applied to all the three arguments.

```
mysql> SELECT 2 BETWEEN 1 AND 3, 2 BETWEEN 3 and 1;
        -> 1, 0
mysql> SELECT 1 BETWEEN 2 AND 3;
        -> 0
mysql> SELECT 'b' BETWEEN 'a' AND 'c';
        -> 1
mysql> SELECT 2 BETWEEN 2 AND '3';
        -> 1
mysql> SELECT 2 BETWEEN 2 AND 'x-3';
        -> 0
```

  For best results when using `BETWEEN` with date or time values, use `CAST()` to explicitly convert the values to the desired data type. Examples: If you compare a `DATETIME` to two `DATE` values, convert the `DATE` values to `DATETIME` values. If you use a string constant such as `'2001-1-1'` in a comparison to a `DATE`, cast the string to a `DATE`.

- *`expr`* `NOT BETWEEN` *`min`* `AND` *`max`*

  This is the same as `NOT (`*`expr`* `BETWEEN` *`min`* `AND` *`max`*`)`.

- `COALESCE(`*`value`*`,...)`

  Returns the first non-`NULL` value in the list, or `NULL` if there are no non-`NULL` values.

  The return type of `COALESCE()` is the aggregated type of the argument types.

```
mysql> SELECT COALESCE(NULL,1);
        -> 1
mysql> SELECT COALESCE(NULL,NULL,NULL);
        -> NULL
```

- `GREATEST(`*`value1`*`,`*`value2`*`,...)`

  With two or more arguments, returns the largest (maximum-valued) argument. The arguments are compared using the same rules as for `LEAST()`.

```
mysql> SELECT GREATEST(2,0);
        -> 2
mysql> SELECT GREATEST(34.0,3.0,5.0,767.0);
        -> 767.0
mysql> SELECT GREATEST('B','A','C');
        -> 'C'
```

`GREATEST()` returns `NULL` if any argument is `NULL`.

- *`expr`* `IN (`*`value`*`,...)`

  Returns `1` if *`expr`* is equal to any of the values in the `IN` list, else returns `0`. If all values are constants, they are evaluated according to the type of *`expr`* and sorted. The search for the item then is done using a binary search. This means `IN` is very quick if the `IN` value list consists entirely of constants. Otherwise, type conversion takes place according to the rules described in Section 12.2, “Type Conversion in Expression Evaluation”, but applied to all the arguments.

```
mysql> SELECT 2 IN (0,3,5,7);
        -> 0
mysql> SELECT 'wefwf' IN ('wee','wefwf','weg');
        -> 1
```

`IN` can be used to compare row constructors:

```
mysql> SELECT (3,4) IN ((1,2), (3,4));
        -> 1
mysql> SELECT (3,4) IN ((1,2), (3,5));
        -> 0
```

You should never mix quoted and unquoted values in an `IN` list because the comparison rules for quoted values (such as strings) and unquoted values (such as numbers) differ. Mixing types may therefore lead to inconsistent results. For example, do not write an `IN` expression like this:

```
SELECT val1 FROM tbl1 WHERE val1 IN (1,2,'a');
```

Instead, write it like this:

```
SELECT val1 FROM tbl1 WHERE val1 IN ('1','2','a');
```

The number of values in the `IN` list is only limited by the `max_allowed_packet` value.

To comply with the SQL standard, `IN` returns `NULL` not only if the expression on the left hand side is `NULL`, but also if no match is found in the list and one of the expressions in the list is `NULL`.

`IN()` syntax can also be used to write certain types of subqueries. See Section 13.2.11.3, “Subqueries with ANY, IN, or SOME”.

- *`expr`* `NOT IN (`*`value`*`,...)`

  This is the same as `NOT (`*`expr`* `IN (`*`value`*`,...))`.

- `ISNULL(`*`expr`*`)`

  If *`expr`* is `NULL`, `ISNULL()` returns `1`, otherwise it returns `0`.

```
mysql> SELECT ISNULL(1+1);
        -> 0
mysql> SELECT ISNULL(1/0);
        -> 1
```

`ISNULL()` can be used instead of `=` to test whether a value is `NULL`. (Comparing a value to `NULL` using `=` always yields `NULL`.)

The `ISNULL()` function shares some special behaviors with the `IS NULL` comparison operator. See the description of `IS NULL`.

- `INTERVAL(N,N1,N2,N3,...)`

  Returns `0` if *`N`* < *`N1`*, `1` if *`N`* < *`N2`* and so on or `-1` if *`N`* is `NULL`. All arguments are treated as integers. It is required that *`N1`* < *`N2`* < *`N3`* < `...` < *`Nn`* for this function to work correctly. This is because a binary search is used (very fast).

```
mysql> SELECT INTERVAL(23, 1, 15, 17, 30, 44, 200);
        -> 3
mysql> SELECT INTERVAL(10, 1, 10, 100, 1000);
        -> 2
mysql> SELECT INTERVAL(22, 23, 30, 44, 200);
        -> 0
```

- `LEAST(value1,value2,...)`

  With two or more arguments, returns the smallest (minimum-valued) argument. The arguments are compared using the following rules:

  - If any argument is `NULL`, the result is `NULL`. No comparison is needed.
  - If all arguments are integer-valued, they are compared as integers.
  - If at least one argument is double precision, they are compared as double-precision values. Otherwise, if at least one argument is a `DECIMAL` value, they are compared as `DECIMAL` values.
  - If the arguments comprise a mix of numbers and strings, they are compared as numbers.
  - If any argument is a nonbinary (character) string, the arguments are compared as nonbinary strings.
  - In all other cases, the arguments are compared as binary strings.

  The return type of `LEAST()` is the aggregated type of the comparison argument types.

```
mysql> SELECT LEAST(2,0);
        -> 0
mysql> SELECT LEAST(34.0,3.0,5.0,767.0);
        -> 3.0
mysql> SELECT LEAST('B','A','C');
        -> 'A'
```

12.3.3 Logical Operators

Table 12.4 Logical Operators

Name	Description
`AND`, `&&`	Logical AND
`NOT`, `!`	Negates value
`||`, `OR`	Logical OR
`XOR`	Logical XOR

In SQL, all logical operators evaluate to `TRUE`, `FALSE`, or `NULL` (`UNKNOWN`). In MySQL, these are implemented as 1 (`TRUE`), 0 (`FALSE`), and `NULL`. Most of this is common to different SQL database servers, although some servers may return any nonzero value for `TRUE`.

MySQL evaluates any nonzero, non-NULL value to TRUE. For example, the following statements all assess to TRUE:

```
mysql> SELECT 10 IS TRUE;
-> 1
mysql> SELECT -10 IS TRUE;
-> 1
mysql> SELECT 'string' IS NOT NULL;
-> 1
```

- NOT, !

  Logical NOT. Evaluates to 1 if the operand is 0, to 0 if the operand is nonzero, and NOT NULL returns NULL.

  ```
  mysql> SELECT NOT 10;
          -> 0
  mysql> SELECT NOT 0;
          -> 1
  mysql> SELECT NOT NULL;
          -> NULL
  mysql> SELECT ! (1+1);
          -> 0
  mysql> SELECT ! 1+1;
          -> 1
  ```

  The last example produces 1 because the expression evaluates the same way as (!1)+1.

- AND, &&

  Logical AND. Evaluates to 1 if all operands are nonzero and not NULL, to 0 if one or more operands are 0, otherwise NULL is returned.

  ```
  mysql> SELECT 1 AND 1;
          -> 1
  mysql> SELECT 1 AND 0;
          -> 0
  mysql> SELECT 1 AND NULL;
          -> NULL
  mysql> SELECT 0 AND NULL;
          -> 0
  mysql> SELECT NULL AND 0;
          -> 0
  ```

- OR, ||

  Logical OR. When both operands are non-NULL, the result is 1 if any operand is nonzero, and 0 otherwise. With a NULL operand, the result is 1 if the other operand is nonzero, and NULL otherwise. If both operands are NULL, the result is NULL.

  ```
  mysql> SELECT 1 OR 1;
          -> 1
  mysql> SELECT 1 OR 0;
          -> 1
  mysql> SELECT 0 OR 0;
          -> 0
  mysql> SELECT 0 OR NULL;
          -> NULL
  mysql> SELECT 1 OR NULL;
          -> 1
  ```

- XOR

  Logical XOR. Returns NULL if either operand is NULL. For non-NULL operands, evaluates to 1 if an odd number of operands is nonzero, otherwise 0 is returned.

```
mysql> SELECT 1 XOR 1;
        -> 0
mysql> SELECT 1 XOR 0;
        -> 1
mysql> SELECT 1 XOR NULL;
        -> NULL
mysql> SELECT 1 XOR 1 XOR 1;
        -> 1
```

`a XOR b` is mathematically equal to `(a AND (NOT b)) OR ((NOT a) and b)`.

12.3.4 Assignment Operators

Table 12.5 Assignment Operators

Name	Description
`=`	Assign a value (as part of a `SET` statement, or as part of the `SET` clause in an `UPDATE` statement)
`:=`	Assign a value

- `:=`

  Assignment operator. Causes the user variable on the left hand side of the operator to take on the value to its right. The value on the right hand side may be a literal value, another variable storing a value, or any legal expression that yields a scalar value, including the result of a query (provided that this value is a scalar value). You can perform multiple assignments in the same `SET` statement. You can perform multiple assignments in the same statement.

  Unlike `=`, the `:=` operator is never interpreted as a comparison operator. This means you can use `:=` in any valid SQL statement (not just in `SET` statements) to assign a value to a variable.

```
mysql> SELECT @var1, @var2;
        -> NULL, NULL
mysql> SELECT @var1 := 1, @var2;
        -> 1, NULL
mysql> SELECT @var1, @var2;
        -> 1, NULL
mysql> SELECT @var1, @var2 := @var1;
        -> 1, 1
mysql> SELECT @var1, @var2;
        -> 1, 1

mysql> SELECT @var1:=COUNT(*) FROM t1;
        -> 4
mysql> SELECT @var1;
        -> 4
```

  You can make value assignments using `:=` in other statements besides `SELECT`, such as `UPDATE`, as shown here:

```
mysql> SELECT @var1;
        -> 4
mysql> SELECT * FROM t1;
        -> 1, 3, 5, 7

mysql> UPDATE t1 SET c1 = 2 WHERE c1 = @var1:= 1;
Query OK, 1 row affected (0.00 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT @var1;
        -> 1
mysql> SELECT * FROM t1;
```

```
        -> 2, 3, 5, 7
```

While it is also possible both to set and to read the value of the same variable in a single SQL statement using the `:=` operator, this is not recommended. Section 9.4, “User-Defined Variables”, explains why you should avoid doing this.

- `=`

  This operator is used to perform value assignments in two cases, described in the next two paragraphs.

  Within a `SET` statement, `=` is treated as an assignment operator that causes the user variable on the left hand side of the operator to take on the value to its right. (In other words, when used in a `SET` statement, `=` is treated identically to `:=`.) The value on the right hand side may be a literal value, another variable storing a value, or any legal expression that yields a scalar value, including the result of a query (provided that this value is a scalar value). You can perform multiple assignments in the same `SET` statement.

  In the `SET` clause of an `UPDATE` statement, `=` also acts as an assignment operator; in this case, however, it causes the column named on the left hand side of the operator to assume the value given to the right, provided any `WHERE` conditions that are part of the `UPDATE` are met. You can make multiple assignments in the same `SET` clause of an `UPDATE` statement.

  In any other context, `=` is treated as a comparison operator.

```
mysql> SELECT @var1, @var2;
        -> NULL, NULL
mysql> SELECT @var1 := 1, @var2;
        -> 1, NULL
mysql> SELECT @var1, @var2;
        -> 1, NULL
mysql> SELECT @var1, @var2 := @var1;
        -> 1, 1
mysql> SELECT @var1, @var2;
        -> 1, 1
```

For more information, see Section 13.7.4.1, “SET Syntax for Variable Assignment”, Section 13.2.12, “UPDATE Syntax”, and Section 13.2.11, “Subquery Syntax”.

12.4 Control Flow Functions

Table 12.6 Flow Control Operators

Name	Description
`CASE`	Case operator
`IF()`	If/else construct
`IFNULL()`	Null if/else construct
`NULLIF()`	Return NULL if expr1 = expr2

- `CASE` *`value`* `WHEN [`*`compare_value`*`] THEN` *`result`* `[WHEN [`*`compare_value`*`] THEN` *`result ...`*`] [ELSE` *`result`*`] END`

  `CASE WHEN [`*`condition`*`] THEN` *`result`* `[WHEN [`*`condition`*`] THEN` *`result ...`*`] [ELSE` *`result`*`] END`

  The first `CASE` syntax returns the *`result`* for the first *`value=compare_value`* comparison that is true. The second syntax returns the result for the first condition that is true. If no comparison or condition is true, the result after `ELSE` is returned, or `NULL` if there is no `ELSE` part.

Note

The syntax of the `CASE` *expression* described here differs slightly from that of the SQL `CASE` *statement* described in Section 13.6.5.1, “CASE Syntax”, for use inside stored programs. The `CASE` statement cannot have an `ELSE NULL` clause, and it is terminated with `END CASE` instead of `END`.

The return type of a `CASE` expression result is the aggregated type of all result values:

- If all types are numeric, the aggregated type is also numeric:
  - If at least one argument is double precision, the result is double precision.
  - Otherwise, if at least one argument is `DECIMAL`, the result is `DECIMAL`.
  - Otherwise, the result is an integer type (with one exception):
    - If all integer types are all signed or all unsigned, the result is the same sign and the precision is the highest of all specified integer types (that is, `TINYINT`, `SMALLINT`, `MEDIUMINT`, `INT`, or `BIGINT`).
    - If there is a combination of signed and unsigned integer types, the result is signed and the precision may be higher. For example, if the types are signed `INT` and unsigned `INT`, the result is signed `BIGINT`.
    - The exception is unsigned `BIGINT` combined with any signed integer type. The result is `DECIMAL` with sufficient precision and scale 0.
- If all types are `BIT`, the result is `BIT`. Otherwise, `BIT` arguments are treated similar to `BIGINT`.
- If all types are `YEAR`, the result is `YEAR`. Otherwise, `YEAR` arguments are treated similar to `INT`.
- If all types are character string (`CHAR` or `VARCHAR`), the result is `VARCHAR` with maximum length determined by the longest character length of the operands.
- If all types are character or binary string, the result is `VARBINARY`.
- `SET` and `ENUM` are treated similar to `VARCHAR`; the result is `VARCHAR`.
- If all types are `JSON`, the result is `JSON`.
- If all types are temporal, the result is temporal:
  - If all temporal types are `DATE`, `TIME`, or `TIMESTAMP`, the result is `DATE`, `TIME`, or `TIMESTAMP`, respectively.
  - Otherwise, for a mix of temporal types, the result is `DATETIME`.
- If all types are `GEOMETRY`, the result is `GEOMETRY`.
- If any type is `BLOB`, the result is `BLOB`.
- For all other type combinations, the result is `VARCHAR`.
- Literal `NULL` operands are ignored for type aggregation.

```
mysql> SELECT CASE 1 WHEN 1 THEN 'one'
    ->     WHEN 2 THEN 'two' ELSE 'more' END;
        -> 'one'
mysql> SELECT CASE WHEN 1>0 THEN 'true' ELSE 'false' END;
        -> 'true'
```

```
mysql> SELECT CASE BINARY 'B'
    ->     WHEN 'a' THEN 1 WHEN 'b' THEN 2 END;
        -> NULL
```

- `IF(expr1,expr2,expr3)`

  If *`expr1`* is `TRUE` (*`expr1 <> 0`* and *`expr1 <> NULL`*), `IF()` returns *`expr2`*. Otherwise, it returns *`expr3`*.

  Note

  There is also an `IF` *statement*, which differs from the `IF()` *function* described here. See Section 13.6.5.2, “IF Syntax”.

  If only one of *`expr2`* or *`expr3`* is explicitly `NULL`, the result type of the `IF()` function is the type of the non-`NULL` expression.

  The default return type of `IF()` (which may matter when it is stored into a temporary table) is calculated as follows:

  - If *`expr2`* or *`expr3`* produce a string, the result is a string.

    If *`expr2`* and *`expr3`* are both strings, the result is case sensitive if either string is case sensitive.

  - If *`expr2`* or *`expr3`* produce a floating-point value, the result is a floating-point value.

  - If *`expr2`* or *`expr3`* produce an integer, the result is an integer.

```
mysql> SELECT IF(1>2,2,3);
        -> 3
mysql> SELECT IF(1<2,'yes','no');
        -> 'yes'
mysql> SELECT IF(STRCMP('test','test1'),'no','yes');
        -> 'no'
```

- `IFNULL(expr1,expr2)`

  If *`expr1`* is not `NULL`, `IFNULL()` returns *`expr1`*; otherwise it returns *`expr2`*.

```
mysql> SELECT IFNULL(1,0);
        -> 1
mysql> SELECT IFNULL(NULL,10);
        -> 10
mysql> SELECT IFNULL(1/0,10);
        -> 10
mysql> SELECT IFNULL(1/0,'yes');
        -> 'yes'
```

  The default return type of `IFNULL(expr1,expr2)` is the more “general” of the two expressions, in the order `STRING`, `REAL`, or `INTEGER`. Consider the case of a table based on expressions or where MySQL must internally store a value returned by `IFNULL()` in a temporary table:

```
mysql> CREATE TABLE tmp SELECT IFNULL(1,'test') AS test;
mysql> DESCRIBE tmp;
+-------+--------------+------+-----+---------+-------+
| Field | Type         | Null | Key | Default | Extra |
+-------+--------------+------+-----+---------+-------+
| test  | varbinary(4) | NO   |     |         |       |
+-------+--------------+------+-----+---------+-------+
```

  In this example, the type of the `test` column is `VARBINARY(4)` (a string type).

- `NULLIF(expr1,expr2)`

Returns `NULL` if *`expr1 = expr2`* is true, otherwise returns *`expr1`*. This is the same as `CASE WHEN` *`expr1 = expr2`* `THEN NULL ELSE` *`expr1`* `END`.

The return value has the same type as the first argument.

```
mysql> SELECT NULLIF(1,1);
        -> NULL
mysql> SELECT NULLIF(1,2);
        -> 1
```

Note

MySQL evaluates *`expr1`* twice if the arguments are not equal.

12.5 String Functions

Table 12.7 String Operators

Name	Description
`ASCII()`	Return numeric value of left-most character
`BIN()`	Return a string containing binary representation of a number
`BIT_LENGTH()`	Return length of argument in bits
`CHAR()`	Return the character for each integer passed
`CHAR_LENGTH()`	Return number of characters in argument
`CHARACTER_LENGTH()`	Synonym for CHAR_LENGTH()
`CONCAT()`	Return concatenated string
`CONCAT_WS()`	Return concatenate with separator
`ELT()`	Return string at index number
`EXPORT_SET()`	Return a string such that for every bit set in the value bits, you get an on string and for every unset bit, you get an off string
`FIELD()`	Return the index (position) of the first argument in the subsequent arguments
`FIND_IN_SET()`	Return the index position of the first argument within the second argument
`FORMAT()`	Return a number formatted to specified number of decimal places
`FROM_BASE64()`	Decode to a base-64 string and return result
`HEX()`	Return a hexadecimal representation of a decimal or string value
`INSERT()`	Insert a substring at the specified position up to the specified number of characters
`INSTR()`	Return the index of the first occurrence of substring
`LCASE()`	Synonym for LOWER()
`LEFT()`	Return the leftmost number of characters as specified
`LENGTH()`	Return the length of a string in bytes
`LIKE`	Simple pattern matching
`LOAD_FILE()`	Load the named file
`LOCATE()`	Return the position of the first occurrence of substring

Name	Description
LOWER()	Return the argument in lowercase
LPAD()	Return the string argument, left-padded with the specified string
LTRIM()	Remove leading spaces
MAKE_SET()	Return a set of comma-separated strings that have the corresponding bit in bits set
MATCH	Perform full-text search
MID()	Return a substring starting from the specified position
NOT LIKE	Negation of simple pattern matching
NOT REGEXP	Negation of REGEXP
OCT()	Return a string containing octal representation of a number
OCTET_LENGTH()	Synonym for LENGTH()
ORD()	Return character code for leftmost character of the argument
POSITION()	Synonym for LOCATE()
QUOTE()	Escape the argument for use in an SQL statement
REGEXP	Pattern matching using regular expressions
REPEAT()	Repeat a string the specified number of times
REPLACE()	Replace occurrences of a specified string
REVERSE()	Reverse the characters in a string
RIGHT()	Return the specified rightmost number of characters
RLIKE	Synonym for REGEXP
RPAD()	Append string the specified number of times
RTRIM()	Remove trailing spaces
SOUNDEX()	Return a soundex string
SOUNDS LIKE	Compare sounds
SPACE()	Return a string of the specified number of spaces
STRCMP()	Compare two strings
SUBSTR()	Return the substring as specified
SUBSTRING()	Return the substring as specified
SUBSTRING_INDEX()	Return a substring from a string before the specified number of occurrences of the delimiter
TO_BASE64()	Return the argument converted to a base-64 string
TRIM()	Remove leading and trailing spaces
UCASE()	Synonym for UPPER()
UNHEX()	Return a string containing hex representation of a number
UPPER()	Convert to uppercase
WEIGHT_STRING()	Return the weight string for a string

String-valued functions return `NULL` if the length of the result would be greater than the value of the `max_allowed_packet` system variable. See Section 5.1.1, “Configuring the Server”.

For functions that operate on string positions, the first position is numbered 1.

For functions that take length arguments, noninteger arguments are rounded to the nearest integer.

- `ASCII(str)`

  Returns the numeric value of the leftmost character of the string *`str`*. Returns `0` if *`str`* is the empty string. Returns `NULL` if *`str`* is `NULL`. `ASCII()` works for 8-bit characters.

```
mysql> SELECT ASCII('2');
        -> 50
mysql> SELECT ASCII(2);
        -> 50
mysql> SELECT ASCII('dx');
        -> 100
```

  See also the `ORD()` function.

- `BIN(N)`

  Returns a string representation of the binary value of *`N`*, where *`N`* is a longlong (`BIGINT`) number. This is equivalent to `CONV(N,10,2)`. Returns `NULL` if *`N`* is `NULL`.

```
mysql> SELECT BIN(12);
        -> '1100'
```

- `BIT_LENGTH(str)`

  Returns the length of the string *`str`* in bits.

```
mysql> SELECT BIT_LENGTH('text');
        -> 32
```

- `CHAR(N,... [USING charset_name])`

  `CHAR()` interprets each argument *`N`* as an integer and returns a string consisting of the characters given by the code values of those integers. `NULL` values are skipped.

```
mysql> SELECT CHAR(77,121,83,81,'76');
        -> 'MySQL'
mysql> SELECT CHAR(77,77.3,'77.3');
        -> 'MMM'
```

  `CHAR()` arguments larger than 255 are converted into multiple result bytes. For example, `CHAR(256)` is equivalent to `CHAR(1,0)`, and `CHAR(256*256)` is equivalent to `CHAR(1,0,0)`:

```
mysql> SELECT HEX(CHAR(1,0)), HEX(CHAR(256));
+----------------+----------------+
| HEX(CHAR(1,0)) | HEX(CHAR(256)) |
+----------------+----------------+
| 0100           | 0100           |
+----------------+----------------+
mysql> SELECT HEX(CHAR(1,0,0)), HEX(CHAR(256*256));
+------------------+--------------------+
| HEX(CHAR(1,0,0)) | HEX(CHAR(256*256)) |
+------------------+--------------------+
| 010000           | 010000             |
+------------------+--------------------+
```

  By default, `CHAR()` returns a binary string. To produce a string in a given character set, use the optional `USING` clause:

```
mysql> SELECT CHARSET(CHAR(X'65')), CHARSET(CHAR(X'65' USING utf8));
+----------------------+------------------------------------+
| CHARSET(CHAR(X'65')) | CHARSET(CHAR(X'65' USING utf8)) |
+----------------------+------------------------------------+
```

```
| binary               | utf8                           |
+----------------------+--------------------------------+
```

If USING is given and the result string is illegal for the given character set, a warning is issued. Also, if strict SQL mode is enabled, the result from CHAR() becomes NULL.

- CHAR_LENGTH(*str*)

  Returns the length of the string *str*, measured in characters. A multibyte character counts as a single character. This means that for a string containing five 2-byte characters, LENGTH() returns 10, whereas CHAR_LENGTH() returns 5.

- CHARACTER_LENGTH(*str*)

  CHARACTER_LENGTH() is a synonym for CHAR_LENGTH().

- CONCAT(*str1*,*str2*,...)

  Returns the string that results from concatenating the arguments. May have one or more arguments. If all arguments are nonbinary strings, the result is a nonbinary string. If the arguments include any binary strings, the result is a binary string. A numeric argument is converted to its equivalent nonbinary string form.

  CONCAT() returns NULL if any argument is NULL.

```
mysql> SELECT CONCAT('My', 'S', 'QL');
        -> 'MySQL'
mysql> SELECT CONCAT('My', NULL, 'QL');
        -> NULL
mysql> SELECT CONCAT(14.3);
        -> '14.3'
```

  For quoted strings, concatenation can be performed by placing the strings next to each other:

```
mysql> SELECT 'My' 'S' 'QL';
        -> 'MySQL'
```

- CONCAT_WS(*separator*,*str1*,*str2*,...)

  CONCAT_WS() stands for Concatenate With Separator and is a special form of CONCAT(). The first argument is the separator for the rest of the arguments. The separator is added between the strings to be concatenated. The separator can be a string, as can the rest of the arguments. If the separator is NULL, the result is NULL.

```
mysql> SELECT CONCAT_WS(',','First name','Second name','Last Name');
        -> 'First name,Second name,Last Name'
mysql> SELECT CONCAT_WS(',','First name',NULL,'Last Name');
        -> 'First name,Last Name'
```

  CONCAT_WS() does not skip empty strings. However, it does skip any NULL values after the separator argument.

- ELT(*N*,*str1*,*str2*,*str3*,...)

  ELT() returns the *N*th element of the list of strings: *str1* if *N* = 1, *str2* if *N* = 2, and so on. Returns NULL if *N* is less than 1 or greater than the number of arguments. ELT() is the complement of FIELD().

```
mysql> SELECT ELT(1, 'ej', 'Heja', 'hej', 'foo');
        -> 'ej'
mysql> SELECT ELT(4, 'ej', 'Heja', 'hej', 'foo');
        -> 'foo'
```

- `EXPORT_SET(bits,on,off[,separator[,number_of_bits]])`

  Returns a string such that for every bit set in the value *`bits`*, you get an *`on`* string and for every bit not set in the value, you get an *`off`* string. Bits in *`bits`* are examined from right to left (from low-order to high-order bits). Strings are added to the result from left to right, separated by the *`separator`* string (the default being the comma character `,`). The number of bits examined is given by *`number_of_bits`*, which has a default of 64 if not specified. *`number_of_bits`* is silently clipped to 64 if larger than 64. It is treated as an unsigned integer, so a value of −1 is effectively the same as 64.

```
mysql> SELECT EXPORT_SET(5,'Y','N',',',4);
        -> 'Y,N,Y,N'
mysql> SELECT EXPORT_SET(6,'1','0',',',10);
        -> '0,1,1,0,0,0,0,0,0,0'
```

- `FIELD(str,str1,str2,str3,...)`

  Returns the index (position) of *`str`* in the *`str1`*, *`str2`*, *`str3`*, `...` list. Returns `0` if *`str`* is not found.

  If all arguments to `FIELD()` are strings, all arguments are compared as strings. If all arguments are numbers, they are compared as numbers. Otherwise, the arguments are compared as double.

  If *`str`* is `NULL`, the return value is `0` because `NULL` fails equality comparison with any value. `FIELD()` is the complement of `ELT()`.

```
mysql> SELECT FIELD('ej', 'Hej', 'ej', 'Heja', 'hej', 'foo');
        -> 2
mysql> SELECT FIELD('fo', 'Hej', 'ej', 'Heja', 'hej', 'foo');
        -> 0
```

- `FIND_IN_SET(str,strlist)`

  Returns a value in the range of 1 to *`N`* if the string *`str`* is in the string list *`strlist`* consisting of *`N`* substrings. A string list is a string composed of substrings separated by `,` characters. If the first argument is a constant string and the second is a column of type `SET`, the `FIND_IN_SET()` function is optimized to use bit arithmetic. Returns `0` if *`str`* is not in *`strlist`* or if *`strlist`* is the empty string. Returns `NULL` if either argument is `NULL`. This function does not work properly if the first argument contains a comma (`,`) character.

```
mysql> SELECT FIND_IN_SET('b','a,b,c,d');
        -> 2
```

- `FORMAT(X,D[,locale])`

  Formats the number *`X`* to a format like `'#,###,###.##'`, rounded to *`D`* decimal places, and returns the result as a string. If *`D`* is `0`, the result has no decimal point or fractional part.

  The optional third parameter enables a locale to be specified to be used for the result number's decimal point, thousands separator, and grouping between separators. Permissible locale values are the same as the legal values for the `lc_time_names` system variable (see Section 10.7, “MySQL Server Locale Support”). If no locale is specified, the default is `'en_US'`.

```
mysql> SELECT FORMAT(12332.123456, 4);
        -> '12,332.1235'
mysql> SELECT FORMAT(12332.1,4);
        -> '12,332.1000'
mysql> SELECT FORMAT(12332.2,0);
        -> '12,332'
mysql> SELECT FORMAT(12332.2,2,'de_DE');
        -> '12.332,20'
```

- `FROM_BASE64(str)`

  Takes a string encoded with the base-64 encoded rules used by `TO_BASE64()` and returns the decoded result as a binary string. The result is `NULL` if the argument is `NULL` or not a valid base-64 string. See the description of `TO_BASE64()` for details about the encoding and decoding rules.

```
mysql> SELECT TO_BASE64('abc'), FROM_BASE64(TO_BASE64('abc'));
       -> 'JWJj', 'abc'
```

- `HEX(str)`, `HEX(N)`

  For a string argument *`str`*, `HEX()` returns a hexadecimal string representation of *`str`* where each byte of each character in *`str`* is converted to two hexadecimal digits. (Multibyte characters therefore become more than two digits.) The inverse of this operation is performed by the `UNHEX()` function.

  For a numeric argument *`N`*, `HEX()` returns a hexadecimal string representation of the value of *`N`* treated as a longlong (`BIGINT`) number. This is equivalent to `CONV(N,10,16)`. The inverse of this operation is performed by `CONV(HEX(N),16,10)`.

```
mysql> SELECT X'616263', HEX('abc'), UNHEX(HEX('abc'));
       -> 'abc', 616263, 'abc'
mysql> SELECT HEX(255), CONV(HEX(255),16,10);
       -> 'FF', 255
```

- `INSERT(str,pos,len,newstr)`

  Returns the string *`str`*, with the substring beginning at position *`pos`* and *`len`* characters long replaced by the string *`newstr`*. Returns the original string if *`pos`* is not within the length of the string. Replaces the rest of the string from position *`pos`* if *`len`* is not within the length of the rest of the string. Returns `NULL` if any argument is `NULL`.

```
mysql> SELECT INSERT('Quadratic', 3, 4, 'What');
       -> 'QuWhattic'
mysql> SELECT INSERT('Quadratic', -1, 4, 'What');
       -> 'Quadratic'
mysql> SELECT INSERT('Quadratic', 3, 100, 'What');
       -> 'QuWhat'
```

  This function is multibyte safe.

- `INSTR(str,substr)`

  Returns the position of the first occurrence of substring *`substr`* in string *`str`*. This is the same as the two-argument form of `LOCATE()`, except that the order of the arguments is reversed.

```
mysql> SELECT INSTR('foobarbar', 'bar');
       -> 4
mysql> SELECT INSTR('xbar', 'foobar');
       -> 0
```

  This function is multibyte safe, and is case sensitive only if at least one argument is a binary string.

- `LCASE(str)`

  `LCASE()` is a synonym for `LOWER()`.

  `LCASE()` used in a view is rewritten as `LOWER()` when storing the view's definition. (Bug #12844279)

- `LEFT(str,len)`

  Returns the leftmost *`len`* characters from the string *`str`*, or `NULL` if any argument is `NULL`.

```
mysql> SELECT LEFT('foobarbar', 5);
        -> 'fooba'
```

This function is multibyte safe.

- `LENGTH(str)`

  Returns the length of the string *`str`*, measured in bytes. A multibyte character counts as multiple bytes. This means that for a string containing five 2-byte characters, `LENGTH()` returns `10`, whereas `CHAR_LENGTH()` returns `5`.

```
mysql> SELECT LENGTH('text');
        -> 4
```

Note

The `Length()` OpenGIS spatial function is named `ST_Length()` in MySQL.

- `LOAD_FILE(file_name)`

  Reads the file and returns the file contents as a string. To use this function, the file must be located on the server host, you must specify the full path name to the file, and you must have the `FILE` privilege. The file must be readable by all and its size less than `max_allowed_packet` bytes. If the `secure_file_priv` system variable is set to a nonempty directory name, the file to be loaded must be located in that directory.

  If the file does not exist or cannot be read because one of the preceding conditions is not satisfied, the function returns `NULL`.

  The `character_set_filesystem` system variable controls interpretation of file names that are given as literal strings.

```
mysql> UPDATE t
           SET blob_col=LOAD_FILE('/tmp/picture')
           WHERE id=1;
```

- `LOCATE(substr,str)`, `LOCATE(substr,str,pos)`

  The first syntax returns the position of the first occurrence of substring *`substr`* in string *`str`*. The second syntax returns the position of the first occurrence of substring *`substr`* in string *`str`*, starting at position *`pos`*. Returns `0` if *`substr`* is not in *`str`*. Returns `NULL` if any argument is `NULL`.

```
mysql> SELECT LOCATE('bar', 'foobarbar');
        -> 4
mysql> SELECT LOCATE('xbar', 'foobar');
        -> 0
mysql> SELECT LOCATE('bar', 'foobarbar', 5);
        -> 7
```

This function is multibyte safe, and is case-sensitive only if at least one argument is a binary string.

- `LOWER(str)`

  Returns the string *`str`* with all characters changed to lowercase according to the current character set mapping. The default is `utf8mb4`.

```
mysql> SELECT LOWER('QUADRATICALLY');
        -> 'quadratically'
```

`LOWER()` (and `UPPER()`) are ineffective when applied to binary strings (`BINARY`, `VARBINARY`, `BLOB`). To perform lettercase conversion, convert the string to a nonbinary string:

```
mysql> SET @str = BINARY 'New York';
mysql> SELECT LOWER(@str), LOWER(CONVERT(@str USING utf8mb4));
+-------------+-----------------------------------+
| LOWER(@str) | LOWER(CONVERT(@str USING utf8mb4)) |
+-------------+-----------------------------------+
| New York    | new york                          |
+-------------+-----------------------------------+
```

For collations of Unicode character sets, `LOWER()` and `UPPER()` work according to the Unicode Collation Algorithm (UCA) version in the collation name, if there is one, and UCA 4.0.0 if no version is specified. For example, `utf8mb4_0900_ai_ci` and `utf8_unicode_520_ci` work according to UCA 9.0.0 and 5.2.0, respectively, whereas `utf8_unicode_ci` works according to UCA 4.0.0. See Section 10.1.10.1, “Unicode Character Sets”.

This function is multibyte safe.

`LCASE()` used within views is rewritten as `LOWER()`.

- `LPAD(str,len,padstr)`

  Returns the string *`str`*, left-padded with the string *`padstr`* to a length of *`len`* characters. If *`str`* is longer than *`len`*, the return value is shortened to *`len`* characters.

  ```
  mysql> SELECT LPAD('hi',4,'??');
         -> '??hi'
  mysql> SELECT LPAD('hi',1,'??');
         -> 'h'
  ```

- `LTRIM(str)`

  Returns the string *`str`* with leading space characters removed.

  ```
  mysql> SELECT LTRIM('  barbar');
         -> 'barbar'
  ```

  This function is multibyte safe.

- `MAKE_SET(bits,str1,str2,...)`

  Returns a set value (a string containing substrings separated by `,` characters) consisting of the strings that have the corresponding bit in *`bits`* set. *`str1`* corresponds to bit 0, *`str2`* to bit 1, and so on. `NULL` values in *`str1`*, *`str2`*, `...` are not appended to the result.

  ```
  mysql> SELECT MAKE_SET(1,'a','b','c');
         -> 'a'
  mysql> SELECT MAKE_SET(1 | 4,'hello','nice','world');
         -> 'hello,world'
  mysql> SELECT MAKE_SET(1 | 4,'hello','nice',NULL,'world');
         -> 'hello'
  mysql> SELECT MAKE_SET(0,'a','b','c');
         -> ''
  ```

- `MID(str,pos,len)`

  `MID(str,pos,len)` is a synonym for `SUBSTRING(str,pos,len)`.

- `OCT(N)`

Returns a string representation of the octal value of *N*, where *N* is a longlong (BIGINT) number. This is equivalent to CONV(*N*,10,8). Returns NULL if *N* is NULL.

```
mysql> SELECT OCT(12);
        -> '14'
```

- OCTET_LENGTH(*str*)

  OCTET_LENGTH() is a synonym for LENGTH().

- ORD(*str*)

  If the leftmost character of the string *str* is a multibyte character, returns the code for that character, calculated from the numeric values of its constituent bytes using this formula:

```
  (1st byte code)
+ (2nd byte code * 256)
+ (3rd byte code * 256^2) ...
```

  If the leftmost character is not a multibyte character, ORD() returns the same value as the ASCII() function.

```
mysql> SELECT ORD('2');
        -> 50
```

- POSITION(*substr* IN *str*)

  POSITION(*substr* IN *str*) is a synonym for LOCATE(*substr*,*str*).

- QUOTE(*str*)

  Quotes a string to produce a result that can be used as a properly escaped data value in an SQL statement. The string is returned enclosed by single quotation marks and with each instance of backslash (\), single quote ('), ASCII NUL, and Control+Z preceded by a backslash. If the argument is NULL, the return value is the word “NULL” without enclosing single quotation marks.

```
mysql> SELECT QUOTE('Don\'t!');
        -> 'Don\'t!'
mysql> SELECT QUOTE(NULL);
        -> NULL
```

  For comparison, see the quoting rules for literal strings and within the C API in Section 9.1.1, “String Literals”, and Section 27.7.7.56, “mysql_real_escape_string_quote()”.

- REPEAT(*str*,*count*)

  Returns a string consisting of the string *str* repeated *count* times. If *count* is less than 1, returns an empty string. Returns NULL if *str* or *count* are NULL.

```
mysql> SELECT REPEAT('MySQL', 3);
        -> 'MySQLMySQLMySQL'
```

- REPLACE(*str*,*from_str*,*to_str*)

  Returns the string *str* with all occurrences of the string *from_str* replaced by the string *to_str*. REPLACE() performs a case-sensitive match when searching for *from_str*.

```
mysql> SELECT REPLACE('www.mysql.com', 'w', 'Ww');
        -> 'WwWwWw.mysql.com'
```

This function is multibyte safe.

- `REVERSE(str)`

  Returns the string *`str`* with the order of the characters reversed.

```
mysql> SELECT REVERSE('abc');
        -> 'cba'
```

  This function is multibyte safe.

- `RIGHT(str,len)`

  Returns the rightmost *`len`* characters from the string *`str`*, or `NULL` if any argument is `NULL`.

```
mysql> SELECT RIGHT('foobarbar', 4);
        -> 'rbar'
```

  This function is multibyte safe.

- `RPAD(str,len,padstr)`

  Returns the string *`str`*, right-padded with the string *`padstr`* to a length of *`len`* characters. If *`str`* is longer than *`len`*, the return value is shortened to *`len`* characters.

```
mysql> SELECT RPAD('hi',5,'?');
        -> 'hi???'
mysql> SELECT RPAD('hi',1,'?');
        -> 'h'
```

  This function is multibyte safe.

- `RTRIM(str)`

  Returns the string *`str`* with trailing space characters removed.

```
mysql> SELECT RTRIM('barbar   ');
        -> 'barbar'
```

  This function is multibyte safe.

- `SOUNDEX(str)`

  Returns a soundex string from *`str`*. Two strings that sound almost the same should have identical soundex strings. A standard soundex string is four characters long, but the `SOUNDEX()` function returns an arbitrarily long string. You can use `SUBSTRING()` on the result to get a standard soundex string. All nonalphabetic characters in *`str`* are ignored. All international alphabetic characters outside the A-Z range are treated as vowels.

  Important

  When using `SOUNDEX()`, you should be aware of the following limitations:

  - This function, as currently implemented, is intended to work well with strings that are in the English language only. Strings in other languages may not produce reliable results.
  - This function is not guaranteed to provide consistent results with strings that use multibyte character sets, including `utf-8`.

    We hope to remove these limitations in a future release. See Bug #22638 for more information.

```
mysql> SELECT SOUNDEX('Hello');
        -> 'H400'
mysql> SELECT SOUNDEX('Quadratically');
        -> 'Q36324'
```

Note

This function implements the original Soundex algorithm, not the more popular enhanced version (also described by D. Knuth). The difference is that original version discards vowels first and duplicates second, whereas the enhanced version discards duplicates first and vowels second.

- *expr1* SOUNDS LIKE *expr2*

  This is the same as SOUNDEX(*expr1*) = SOUNDEX(*expr2*).

- SPACE(*N*)

  Returns a string consisting of *N* space characters.

```
mysql> SELECT SPACE(6);
        -> '      '
```

- SUBSTR(*str*,*pos*), SUBSTR(*str* FROM *pos*), SUBSTR(*str*,*pos*,*len*), SUBSTR(*str* FROM *pos* FOR *len*)

  SUBSTR() is a synonym for SUBSTRING().

- SUBSTRING(*str*,*pos*), SUBSTRING(*str* FROM *pos*), SUBSTRING(*str*,*pos*,*len*), SUBSTRING(*str* FROM *pos* FOR *len*)

  The forms without a *len* argument return a substring from string *str* starting at position *pos*. The forms with a *len* argument return a substring *len* characters long from string *str*, starting at position *pos*. The forms that use FROM are standard SQL syntax. It is also possible to use a negative value for *pos*. In this case, the beginning of the substring is *pos* characters from the end of the string, rather than the beginning. A negative value may be used for *pos* in any of the forms of this function.

  For all forms of SUBSTRING(), the position of the first character in the string from which the substring is to be extracted is reckoned as 1.

```
mysql> SELECT SUBSTRING('Quadratically',5);
        -> 'ratically'
mysql> SELECT SUBSTRING('foobarbar' FROM 4);
        -> 'barbar'
mysql> SELECT SUBSTRING('Quadratically',5,6);
        -> 'ratica'
mysql> SELECT SUBSTRING('Sakila', -3);
        -> 'ila'
mysql> SELECT SUBSTRING('Sakila', -5, 3);
        -> 'aki'
mysql> SELECT SUBSTRING('Sakila' FROM -4 FOR 2);
        -> 'ki'
```

  This function is multibyte safe.

  If *len* is less than 1, the result is the empty string.

- SUBSTRING_INDEX(*str*,*delim*,*count*)

  Returns the substring from string *str* before *count* occurrences of the delimiter *delim*. If *count* is positive, everything to the left of the final delimiter (counting from the left) is returned. If *count*

is negative, everything to the right of the final delimiter (counting from the right) is returned. `SUBSTRING_INDEX()` performs a case-sensitive match when searching for *`delim`*.

```
mysql> SELECT SUBSTRING_INDEX('www.mysql.com', '.', 2);
        -> 'www.mysql'
mysql> SELECT SUBSTRING_INDEX('www.mysql.com', '.', -2);
        -> 'mysql.com'
```

This function is multibyte safe.

- `TO_BASE64(`*`str`*`)`

  Converts the string argument to base-64 encoded form and returns the result as a character string with the connection character set and collation. If the argument is not a string, it is converted to a string before conversion takes place. The result is `NULL` if the argument is `NULL`. Base-64 encoded strings can be decoded using the `FROM_BASE64()` function.

  ```
  mysql> SELECT TO_BASE64('abc'), FROM_BASE64(TO_BASE64('abc'));
          -> 'JWJj', 'abc'
  ```

  Different base-64 encoding schemes exist. These are the encoding and decoding rules used by `TO_BASE64()` and `FROM_BASE64()`:

  - The encoding for alphabet value 62 is `'+'`.
  - The encoding for alphabet value 63 is `'/'`.
  - Encoded output consists of groups of 4 printable characters. Each 3 bytes of the input data are encoded using 4 characters. If the last group is incomplete, it is padded with `'='` characters to a length of 4.
  - A newline is added after each 76 characters of encoded output to divide long output into multiple lines.
  - Decoding recognizes and ignores newline, carriage return, tab, and space.

- `TRIM([{BOTH | LEADING | TRAILING} [`*`remstr`*`] FROM] `*`str`*`)`, `TRIM([`*`remstr`*` FROM] `*`str`*`)`

  Returns the string *`str`* with all *`remstr`* prefixes or suffixes removed. If none of the specifiers `BOTH`, `LEADING`, or `TRAILING` is given, `BOTH` is assumed. *`remstr`* is optional and, if not specified, spaces are removed.

  ```
  mysql> SELECT TRIM('  bar   ');
          -> 'bar'
  mysql> SELECT TRIM(LEADING 'x' FROM 'xxxbarxxx');
          -> 'barxxx'
  mysql> SELECT TRIM(BOTH 'x' FROM 'xxxbarxxx');
          -> 'bar'
  mysql> SELECT TRIM(TRAILING 'xyz' FROM 'barxxyz');
          -> 'barx'
  ```

  This function is multibyte safe.

- `UCASE(`*`str`*`)`

  `UCASE()` is a synonym for `UPPER()`.

  `UCASE()` used within views is rewritten as `UPPER()`.

- `UNHEX(`*`str`*`)`

For a string argument *str*, UNHEX(*str*) interprets each pair of characters in the argument as a hexadecimal number and converts it to the byte represented by the number. The return value is a binary string.

```
mysql> SELECT UNHEX('4D7953514C');
        -> 'MySQL'
mysql> SELECT X'4D7953514C';
        -> 'MySQL'
mysql> SELECT UNHEX(HEX('string'));
        -> 'string'
mysql> SELECT HEX(UNHEX('1267'));
        -> '1267'
```

The characters in the argument string must be legal hexadecimal digits: '0' .. '9', 'A' .. 'F', 'a' .. 'f'. If the argument contains any nonhexadecimal digits, the result is NULL:

```
mysql> SELECT UNHEX('GG');
+-------------+
| UNHEX('GG') |
+-------------+
| NULL        |
+-------------+
```

A NULL result can occur if the argument to UNHEX() is a BINARY column, because values are padded with 0x00 bytes when stored but those bytes are not stripped on retrieval. For example, '41' is stored into a CHAR(3) column as '41 ' and retrieved as '41' (with the trailing pad space stripped), so UNHEX() for the column value returns 'A'. By contrast '41' is stored into a BINARY(3) column as '41\0' and retrieved as '41\0' (with the trailing pad 0x00 byte not stripped). '\0' is not a legal hexadecimal digit, so UNHEX() for the column value returns NULL.

For a numeric argument *N*, the inverse of HEX(*N*) is not performed by UNHEX(). Use CONV(HEX(*N*),16,10) instead. See the description of HEX().

- UPPER(*str*)

  Returns the string *str* with all characters changed to uppercase according to the current character set mapping. The default is utf8mb4.

  ```
  mysql> SELECT UPPER('Hej');
          -> 'HEJ'
  ```

  See the description of LOWER() for information that also applies to UPPER(). This included information about how to perform lettercase conversion of binary strings (BINARY, VARBINARY, BLOB) for which these functions are ineffective, and information about case folding for Unicode character sets.

  This function is multibyte safe.

  UCASE() used within views is rewritten as UPPER().

- WEIGHT_STRING(*str* [AS {CHAR|BINARY}(*N*)] [*flags*])

  This function returns the weight string for the input string. The return value is a binary string that represents the comparison and sorting value of the string. It has these properties:

  - If WEIGHT_STRING(*str1*) = WEIGHT_STRING(*str2*), then *str1* = *str2* (*str1* and *str2* are considered equal)
  - If WEIGHT_STRING(*str1*) < WEIGHT_STRING(*str2*), then *str1* < *str2* (*str1* sorts before *str2*)

`WEIGHT_STRING()` is a debugging function intended for internal use. Its behavior can change without notice between MySQL versions. It can be used for testing and debugging of collations, especially if you are adding a new collation. See Section 10.4, “Adding a Collation to a Character Set”.

This list briefly summarizes the arguments. More details are given in the discussion following the list.

- *`str`*: The input string expression.
- `AS` clause: Optional; cast the input string to a given type and length.
- *`flags`*: Optional; unused.

The input string, *`str`*, is a string expression. If the input is a nonbinary (character) string such as a `CHAR`, `VARCHAR`, or `TEXT` value, the return value contains the collation weights for the string. If the input is a binary (byte) string such as a `BINARY`, `VARBINARY`, or `BLOB` value, the return value is the same as the input (the weight for each byte in a binary string is the byte value). If the input is `NULL`, `WEIGHT_STRING()` returns `NULL`.

Examples:

```
mysql> SET @s = _utf8mb4 'AB' COLLATE utf8mb4_0900_ai_ci;
mysql> SELECT @s, HEX(@s), HEX(WEIGHT_STRING(@s));
+------+---------+------------------------+
| @s   | HEX(@s) | HEX(WEIGHT_STRING(@s)) |
+------+---------+------------------------+
| AB   | 4142    | 1C471C60               |
+------+---------+------------------------+
```

```
mysql> SET @s = _utf8mb4 'ab' COLLATE utf8mb4_0900_ai_ci;
mysql> SELECT @s, HEX(@s), HEX(WEIGHT_STRING(@s));
+------+---------+------------------------+
| @s   | HEX(@s) | HEX(WEIGHT_STRING(@s)) |
+------+---------+------------------------+
| ab   | 6162    | 1C471C60               |
+------+---------+------------------------+
```

```
mysql> SET @s = CAST('AB' AS BINARY);
mysql> SELECT @s, HEX(@s), HEX(WEIGHT_STRING(@s));
+------+---------+------------------------+
| @s   | HEX(@s) | HEX(WEIGHT_STRING(@s)) |
+------+---------+------------------------+
| AB   | 4142    | 4142                   |
+------+---------+------------------------+
```

```
mysql> SET @s = CAST('ab' AS BINARY);
mysql> SELECT @s, HEX(@s), HEX(WEIGHT_STRING(@s));
+------+---------+------------------------+
| @s   | HEX(@s) | HEX(WEIGHT_STRING(@s)) |
+------+---------+------------------------+
| ab   | 6162    | 6162                   |
+------+---------+------------------------+
```

The preceding examples use `HEX()` to display the `WEIGHT_STRING()` result. Because the result is a binary value, `HEX()` can be especially useful when the result contains nonprinting values, to display it in printable form:

```
mysql> SET @s = CONVERT(X'C39F' USING utf8) COLLATE utf8_czech_ci;
mysql> SELECT HEX(WEIGHT_STRING(@s));
+------------------------+
| HEX(WEIGHT_STRING(@s)) |
+------------------------+
```

```
| 0FEA0FEA              |
+-----------------------+
```

For non-NULL return values, the data type of the value is VARBINARY if its length is within the maximum length for VARBINARY, otherwise the data type is BLOB.

The AS clause may be given to cast the input string to a nonbinary or binary string and to force it to a given length:

- AS CHAR(*N*) casts the string to a nonbinary string and pads it on the right with spaces to a length of *N* characters. *N* must be at least 1. If *N* is less than the length of the input string, the string is truncated to *N* characters. No warning occurs for truncation.
- AS BINARY(*N*) is similar but casts the string to a binary string, *N* is measured in bytes (not characters), and padding uses 0x00 bytes (not spaces).

```
mysql> SET NAMES 'latin1';
mysql> SELECT HEX(WEIGHT_STRING('ab' AS CHAR(4)));
+-------------------------------------+
| HEX(WEIGHT_STRING('ab' AS CHAR(4))) |
+-------------------------------------+
| 41422020                            |
+-------------------------------------+
mysql> SET NAMES 'utf8';
mysql> SELECT HEX(WEIGHT_STRING('ab' AS CHAR(4)));
+-------------------------------------+
| HEX(WEIGHT_STRING('ab' AS CHAR(4))) |
+-------------------------------------+
| 0041004200200020                    |
+-------------------------------------+
```

```
mysql> SELECT HEX(WEIGHT_STRING('ab' AS BINARY(4)));
+---------------------------------------+
| HEX(WEIGHT_STRING('ab' AS BINARY(4))) |
+---------------------------------------+
| 61620000                              |
+---------------------------------------+
```

The *flags* clause currently is unused.

12.5.1 String Comparison Functions

Table 12.8 String Comparison Operators

Name	Description
LIKE	Simple pattern matching
NOT LIKE	Negation of simple pattern matching
STRCMP()	Compare two strings

If a string function is given a binary string as an argument, the resulting string is also a binary string. A number converted to a string is treated as a binary string. This affects only comparisons.

Normally, if any expression in a string comparison is case sensitive, the comparison is performed in case-sensitive fashion.

- *expr* LIKE *pat* [ESCAPE '*escape_char*']

  Pattern matching using an SQL pattern. Returns 1 (TRUE) or 0 (FALSE). If either *expr* or *pat* is NULL, the result is NULL.

  The pattern need not be a literal string. For example, it can be specified as a string expression or table column.

Per the SQL standard, `LIKE` performs matching on a per-character basis, thus it can produce results different from the `=` comparison operator:

```
mysql> SELECT 'ä' LIKE 'ae' COLLATE latin1_german2_ci;
+-----------------------------------------+
| 'ä' LIKE 'ae' COLLATE latin1_german2_ci |
+-----------------------------------------+
|                                       0 |
+-----------------------------------------+
mysql> SELECT 'ä' = 'ae' COLLATE latin1_german2_ci;
+-------------------------------------+
| 'ä' = 'ae' COLLATE latin1_german2_ci |
+-------------------------------------+
|                                   1 |
+-------------------------------------+
```

In particular, trailing spaces are significant, which is not true for `CHAR` or `VARCHAR` comparisons performed with the `=` operator:

```
mysql> SELECT 'a' = 'a ', 'a' LIKE 'a ';
+------------+---------------+
| 'a' = 'a ' | 'a' LIKE 'a ' |
+------------+---------------+
|          1 |             0 |
+------------+---------------+
1 row in set (0.00 sec)
```

With `LIKE` you can use the following two wildcard characters in the pattern:

- `%` matches any number of characters, even zero characters.
- `_` matches exactly one character.

```
mysql> SELECT 'David!' LIKE 'David_';
        -> 1
mysql> SELECT 'David!' LIKE '%D%v%';
        -> 1
```

To test for literal instances of a wildcard character, precede it by the escape character. If you do not specify the `ESCAPE` character, `\` is assumed.

- `\%` matches one `%` character.
- `\_` matches one `_` character.

```
mysql> SELECT 'David!' LIKE 'David\_';
        -> 0
mysql> SELECT 'David_' LIKE 'David\_';
        -> 1
```

To specify a different escape character, use the `ESCAPE` clause:

```
mysql> SELECT 'David_' LIKE 'David|_' ESCAPE '|';
        -> 1
```

The escape sequence should be empty or one character long. The expression must evaluate as a constant at execution time. If the `NO_BACKSLASH_ESCAPES` SQL mode is enabled, the sequence cannot be empty.

The following two statements illustrate that string comparisons are not case sensitive unless one of the operands is a case sensitive (uses a case-sensitive collation or is a binary string):

```
mysql> SELECT 'abc' LIKE 'ABC';
        -> 1
mysql> SELECT 'abc' LIKE _utf8mb4 'ABC' COLLATE utf8mb4_0900_as_cs;
        -> 0
mysql> SELECT 'abc' LIKE _utf8mb4 'ABC' COLLATE utf8mb4_bin;
        -> 0
mysql> SELECT 'abc' LIKE BINARY 'ABC';
        -> 0
```

As an extension to standard SQL, MySQL permits `LIKE` on numeric expressions.

```
mysql> SELECT 10 LIKE '1%';
        -> 1
```

Note

Because MySQL uses C escape syntax in strings (for example, `\n` to represent a newline character), you must double any `\` that you use in `LIKE` strings. For example, to search for `\n`, specify it as `\\n`. To search for `\`, specify it as `\\\\`; this is because the backslashes are stripped once by the parser and again when the pattern match is made, leaving a single backslash to be matched against.

Exception: At the end of the pattern string, backslash can be specified as `\\`. At the end of the string, backslash stands for itself because there is nothing following to escape. Suppose that a table contains the following values:

```
mysql> SELECT filename FROM t1;
+--------------+
| filename     |
+--------------+
| C:           |
| C:\          |
| C:\Programs  |
| C:\Programs\ |
+--------------+
```

To test for values that end with backslash, you can match the values using either of the following patterns:

```
mysql> SELECT filename, filename LIKE '%\\' FROM t1;
+--------------+---------------------+
| filename     | filename LIKE '%\\' |
+--------------+---------------------+
| C:           |                   0 |
| C:\          |                   1 |
| C:\Programs  |                   0 |
| C:\Programs\ |                   1 |
+--------------+---------------------+

mysql> SELECT filename, filename LIKE '%\\\\' FROM t1;
+--------------+----------------------+
| filename     | filename LIKE '%\\\\' |
+--------------+----------------------+
| C:           |                    0 |
| C:\          |                    1 |
| C:\Programs  |                    0 |
| C:\Programs\ |                    1 |
+--------------+----------------------+
```

- *`expr`* `NOT LIKE` *`pat`* `[ESCAPE '`*`escape_char`*`']`

  This is the same as `NOT (`*`expr`* `LIKE` *`pat`* `[ESCAPE '`*`escape_char`*`'])`.

Note

Aggregate queries involving `NOT LIKE` comparisons with columns containing `NULL` may yield unexpected results. For example, consider the following table and data:

```
CREATE TABLE foo (bar VARCHAR(10));

INSERT INTO foo VALUES (NULL), (NULL);
```

The query `SELECT COUNT(*) FROM foo WHERE bar LIKE '%baz%';` returns `0`. You might assume that `SELECT COUNT(*) FROM foo WHERE bar NOT LIKE '%baz%';` would return `2`. However, this is not the case: The second query returns `0`. This is because `NULL NOT LIKE expr` always returns `NULL`, regardless of the value of *`expr`*. The same is true for aggregate queries involving `NULL` and comparisons using `NOT RLIKE` or `NOT REGEXP`. In such cases, you must test explicitly for `NOT NULL` using `OR` (and not `AND`), as shown here:

```
SELECT COUNT(*) FROM foo WHERE bar NOT LIKE '%baz%' OR bar IS NULL;
```

- `STRCMP(`*`expr1`*`,`*`expr2`*`)`

  `STRCMP()` returns `0` if the strings are the same, `-1` if the first argument is smaller than the second according to the current sort order, and `1` otherwise.

```
mysql> SELECT STRCMP('text', 'text2');
        -> -1
mysql> SELECT STRCMP('text2', 'text');
        -> 1
mysql> SELECT STRCMP('text', 'text');
        -> 0
```

  `STRCMP()` performs the comparison using the collation of the arguments.

```
mysql> SET @s1 = _utf8mb4 'x' COLLATE utf8mb4_0900_ai_ci;
mysql> SET @s2 = _utf8mb4 'X' COLLATE utf8mb4_0900_ai_ci;
mysql> SET @s3 = _utf8mb4 'x' COLLATE utf8mb4_0900_as_cs;
mysql> SET @s4 = _utf8mb4 'X' COLLATE utf8mb4_0900_as_cs;
mysql> SELECT STRCMP(@s1, @s2), STRCMP(@s3, @s4);
+------------------+------------------+
| STRCMP(@s1, @s2) | STRCMP(@s3, @s4) |
+------------------+------------------+
|                0 |               -1 |
+------------------+------------------+
```

  If the collations are incompatible, one of the arguments must be converted to be compatible with the other. See Section 10.1.8.4, “Collation Coercibility in Expressions”.

```
mysql> SET @s1 = _utf8mb4 'x' COLLATE utf8mb4_0900_ai_ci;
mysql> SET @s2 = _utf8mb4 'X' COLLATE utf8mb4_0900_ai_ci;
mysql> SET @s3 = _utf8mb4 'x' COLLATE utf8mb4_0900_as_cs;
mysql> SET @s4 = _utf8mb4 'X' COLLATE utf8mb4_0900_as_cs;
-->
mysql> SELECT STRCMP(@s1, @s3);
ERROR 1267 (HY000): Illegal mix of collations (utf8mb4_0900_ai_ci,IMPLICIT)
and (utf8mb4_0900_as_cs,IMPLICIT) for operation 'strcmp'
mysql> SELECT STRCMP(@s1, @s3 COLLATE utf8mb4_0900_ai_ci);
+--------------------------------------------+
| STRCMP(@s1, @s3 COLLATE utf8mb4_0900_ai_ci) |
+--------------------------------------------+
|                                          0 |
```

```
+-----------------------------------------+
```

12.5.2 Regular Expressions

Table 12.9 String Regular Expression Operators

Name	Description
NOT REGEXP	Negation of REGEXP
REGEXP	Pattern matching using regular expressions
RLIKE	Synonym for REGEXP

A regular expression is a powerful way of specifying a pattern for a complex search.

MySQL uses Henry Spencer's implementation of regular expressions, which is aimed at conformance with POSIX 1003.2. MySQL uses the extended version to support pattern-matching operations performed with the REGEXP operator in SQL statements.

This section summarizes, with examples, the special characters and constructs that can be used in MySQL for REGEXP operations. It does not contain all the details that can be found in Henry Spencer's regex(7) manual page. That manual page is included in MySQL source distributions, in the regex.7 file under the regex directory. See also Section 3.3.4.7, “Pattern Matching”.

Regular Expression Operators

- *expr* NOT REGEXP *pat*, *expr* NOT RLIKE *pat*

  This is the same as NOT (*expr* REGEXP *pat*).

- *expr* REGEXP *pat*, *expr* RLIKE *pat*

  Performs a pattern match of a string expression *expr* against a pattern *pat*. The pattern can be an extended regular expression, the syntax for which is discussed later in this section. Returns 1 if *expr* matches *pat*; otherwise it returns 0. If either *expr* or *pat* is NULL, the result is NULL. RLIKE is a synonym for REGEXP.

  The pattern need not be a literal string. For example, it can be specified as a string expression or table column.

  Note

  Because MySQL uses the C escape syntax in strings (for example, \n to represent the newline character), you must double any \ that you use in your REGEXP strings.

  REGEXP is not case sensitive, except when used with binary strings.

```
mysql> SELECT 'Michael!' REGEXP '.*';
        -> 1
mysql> SELECT 'new*\n*line' REGEXP 'new\\*.\\*line';
        -> 1
mysql> SELECT 'a' REGEXP 'A', 'a' REGEXP BINARY 'A';
        -> 1  0
mysql> SELECT 'a' REGEXP '^[a-d]';
        -> 1
```

  REGEXP and RLIKE use the character set and collations of the arguments when deciding the type of a character and performing the comparison. If the arguments have different character sets or collations, coercibility rules apply as described in Section 10.1.8.4, “Collation Coercibility in Expressions”.

Warning

The `REGEXP` and `RLIKE` operators work in byte-wise fashion, so they are not multibyte safe and may produce unexpected results with multibyte character sets. In addition, these operators compare characters by their byte values and accented characters may not compare as equal even if a given collation treats them as equal.

Syntax of Regular Expressions

A regular expression describes a set of strings. The simplest regular expression is one that has no special characters in it. For example, the regular expression `hello` matches `hello` and nothing else.

Nontrivial regular expressions use certain special constructs so that they can match more than one string. For example, the regular expression `hello|word` matches either the string `hello` or the string `word`.

As a more complex example, the regular expression `B[an]*s` matches any of the strings `Bananas`, `Baaaaas`, `Bs`, and any other string starting with a `B`, ending with an `s`, and containing any number of `a` or `n` characters in between.

A regular expression for the `REGEXP` operator may use any of the following special characters and constructs:

- `^`

  Match the beginning of a string.

```
mysql> SELECT 'fo\nfo' REGEXP '^fo$';                  -> 0
mysql> SELECT 'fofo' REGEXP '^fo';                     -> 1
```

- `$`

  Match the end of a string.

```
mysql> SELECT 'fo\no' REGEXP '^fo\no$';                -> 1
mysql> SELECT 'fo\no' REGEXP '^fo$';                   -> 0
```

- `.`

  Match any character (including carriage return and newline).

```
mysql> SELECT 'fofo' REGEXP '^f.*$';                   -> 1
mysql> SELECT 'fo\r\nfo' REGEXP '^f.*$';               -> 1
```

- `a*`

  Match any sequence of zero or more `a` characters.

```
mysql> SELECT 'Ban' REGEXP '^Ba*n';                    -> 1
mysql> SELECT 'Baaan' REGEXP '^Ba*n';                  -> 1
mysql> SELECT 'Bn' REGEXP '^Ba*n';                     -> 1
```

- `a+`

  Match any sequence of one or more `a` characters.

```
mysql> SELECT 'Ban' REGEXP '^Ba+n';                    -> 1
mysql> SELECT 'Bn' REGEXP '^Ba+n';                     -> 0
```

- `a?`

  Match either zero or one `a` character.

```
mysql> SELECT 'Bn' REGEXP '^Ba?n';                    -> 1
mysql> SELECT 'Ban' REGEXP '^Ba?n';                   -> 1
mysql> SELECT 'Baan' REGEXP '^Ba?n';                  -> 0
```

- `de|abc`

  Match either of the sequences `de` or `abc`.

```
mysql> SELECT 'pi' REGEXP 'pi|apa';                   -> 1
mysql> SELECT 'axe' REGEXP 'pi|apa';                  -> 0
mysql> SELECT 'apa' REGEXP 'pi|apa';                  -> 1
mysql> SELECT 'apa' REGEXP '^(pi|apa)$';              -> 1
mysql> SELECT 'pi' REGEXP '^(pi|apa)$';               -> 1
mysql> SELECT 'pix' REGEXP '^(pi|apa)$';              -> 0
```

- `(abc)*`

  Match zero or more instances of the sequence `abc`.

```
mysql> SELECT 'pi' REGEXP '^(pi)*$';                  -> 1
mysql> SELECT 'pip' REGEXP '^(pi)*$';                 -> 0
mysql> SELECT 'pipi' REGEXP '^(pi)*$';                -> 1
```

- `{1}`, `{2,3}`

  `{n}` or `{m,n}` notation provides a more general way of writing regular expressions that match many occurrences of the previous atom (or “piece”) of the pattern. `m` and `n` are integers.

  - `a*`

    Can be written as `a{0,}`.

  - `a+`

    Can be written as `a{1,}`.

  - `a?`

    Can be written as `a{0,1}`.

  To be more precise, `a{n}` matches exactly `n` instances of `a`. `a{n,}` matches `n` or more instances of `a`. `a{m,n}` matches `m` through `n` instances of `a`, inclusive.

  `m` and `n` must be in the range from `0` to `RE_DUP_MAX` (default 255), inclusive. If both `m` and `n` are given, `m` must be less than or equal to `n`.

```
mysql> SELECT 'abcde' REGEXP 'a[bcd]{2}e';            -> 0
mysql> SELECT 'abcde' REGEXP 'a[bcd]{3}e';            -> 1
mysql> SELECT 'abcde' REGEXP 'a[bcd]{1,10}e';         -> 1
```

- `[a-dX]`, `[^a-dX]`

  Matches any character that is (or is not, if ^ is used) either `a`, `b`, `c`, `d` or `X`. A `-` character between two other characters forms a range that matches all characters from the first character to the second. For example, `[0-9]` matches any decimal digit. To include a literal `]` character, it must immediately follow the opening bracket `[`. To include a literal `-` character, it must be written first or last. Any character that does not have a defined special meaning inside a `[]` pair matches only itself.

```
mysql> SELECT 'aXbc' REGEXP '[a-dXYZ]';                -> 1
mysql> SELECT 'aXbc' REGEXP '^[a-dXYZ]$';              -> 0
mysql> SELECT 'aXbc' REGEXP '^[a-dXYZ]+$';             -> 1
mysql> SELECT 'aXbc' REGEXP '^[^a-dXYZ]+$';            -> 0
mysql> SELECT 'gheis' REGEXP '^[^a-dXYZ]+$';           -> 1
mysql> SELECT 'gheisa' REGEXP '^[^a-dXYZ]+$';          -> 0
```

- `[.characters.]`

  Within a bracket expression (written using `[` and `]`), matches the sequence of characters of that collating element. `characters` is either a single character or a character name like `newline`. The following table lists the permissible character names.

  The following table shows the permissible character names and the characters that they match. For characters given as numeric values, the values are represented in octal.

Name	Character	Name	Character
NUL	0	SOH	001
STX	002	ETX	003
EOT	004	ENQ	005
ACK	006	BEL	007
alert	007	BS	010
backspace	'\b'	HT	011
tab	'\t'	LF	012
newline	'\n'	VT	013
vertical-tab	'\v'	FF	014
form-feed	'\f'	CR	015
carriage-return	'\r'	SO	016
SI	017	DLE	020
DC1	021	DC2	022
DC3	023	DC4	024
NAK	025	SYN	026
ETB	027	CAN	030
EM	031	SUB	032
ESC	033	IS4	034
FS	034	IS3	035
GS	035	IS2	036
RS	036	IS1	037
US	037	space	' '
exclamation-mark	'!'	quotation-mark	'"'
number-sign	'#'	dollar-sign	'$'
percent-sign	'%'	ampersand	'&'
apostrophe	'\''	left-parenthesis	'('
right-parenthesis	')'	asterisk	'*'
plus-sign	'+'	comma	','
hyphen	'-'	hyphen-minus	'-'

Name	Character	Name	Character
period	'.'	full-stop	'.'
slash	'/'	solidus	'/'
zero	'0'	one	'1'
two	'2'	three	'3'
four	'4'	five	'5'
six	'6'	seven	'7'
eight	'8'	nine	'9'
colon	':'	semicolon	';'
less-than-sign	'<'	equals-sign	'='
greater-than-sign	'>'	question-mark	'?'
commercial-at	'@'	left-square-bracket	'['
backslash	'\\'	reverse-solidus	'\\'
right-square-bracket	']'	circumflex	'^'
circumflex-accent	'^'	underscore	'_'
low-line	'_'	grave-accent	'`'
left-brace	'{'	left-curly-bracket	'{'
vertical-line	'|'	right-brace	'}'
right-curly-bracket	'}'	tilde	'~'
DEL	177		

```
mysql> SELECT '~' REGEXP '[[.~.]]';                    -> 1
mysql> SELECT '~' REGEXP '[[.tilde.]]';                 -> 1
```

- `[=character_class=]`

  Within a bracket expression (written using `[` and `]`), `[=character_class=]` represents an equivalence class. It matches all characters with the same collation value, including itself. For example, if `o` and `(+)` are the members of an equivalence class, `[[=o=]]`, `[[=(+)=]]`, and `[o(+)]` are all synonymous. An equivalence class may not be used as an endpoint of a range.

- `[:character_class:]`

  Within a bracket expression (written using `[` and `]`), `[:character_class:]` represents a character class that matches all characters belonging to that class. The following table lists the standard class names. These names stand for the character classes defined in the `ctype(3)` manual page. A particular locale may provide other class names. A character class may not be used as an endpoint of a range.

Character Class Name	Meaning
alnum	Alphanumeric characters
alpha	Alphabetic characters
blank	Whitespace characters
cntrl	Control characters

Character Class Name	Meaning
`digit`	Digit characters
`graph`	Graphic characters
`lower`	Lowercase alphabetic characters
`print`	Graphic or space characters
`punct`	Punctuation characters
`space`	Space, tab, newline, and carriage return
`upper`	Uppercase alphabetic characters
`xdigit`	Hexadecimal digit characters

```
mysql> SELECT 'justalnums' REGEXP '[[:alnum:]]+';       -> 1
mysql> SELECT '!!' REGEXP '[[:alnum:]]+';               -> 0
```

- `[[:<:]]`, `[[:>:]]`

  These markers stand for word boundaries. They match the beginning and end of words, respectively. A word is a sequence of word characters that is not preceded by or followed by word characters. A word character is an alphanumeric character in the `alnum` class or an underscore (_).

```
mysql> SELECT 'a word a' REGEXP '[[:<:]]word[[:>:]]';   -> 1
mysql> SELECT 'a xword a' REGEXP '[[:<:]]word[[:>:]]';  -> 0
```

To use a literal instance of a special character in a regular expression, precede it by two backslash (\) characters. The MySQL parser interprets one of the backslashes, and the regular expression library interprets the other. For example, to match the string `1+2` that contains the special `+` character, only the last of the following regular expressions is the correct one:

```
mysql> SELECT '1+2' REGEXP '1+2';                       -> 0
mysql> SELECT '1+2' REGEXP '1\+2';                      -> 0
mysql> SELECT '1+2' REGEXP '1\\+2';                     -> 1
```

12.5.3 Character Set and Collation of Function Results

MySQL has many operators and functions that return a string. This section answers the question: What is the character set and collation of such a string?

For simple functions that take string input and return a string result as output, the output's character set and collation are the same as those of the principal input value. For example, `UPPER(X)` returns a string with the same character string and collation as *X*. The same applies for `INSTR()`, `LCASE()`, `LOWER()`, `LTRIM()`, `MID()`, `REPEAT()`, `REPLACE()`, `REVERSE()`, `RIGHT()`, `RPAD()`, `RTRIM()`, `SOUNDEX()`, `SUBSTRING()`, `TRIM()`, `UCASE()`, and `UPPER()`.

Note

The `REPLACE()` function, unlike all other functions, always ignores the collation of the string input and performs a case-sensitive comparison.

If a string input or function result is a binary string, the string has the `binary` character set and collation. This can be checked by using the `CHARSET()` and `COLLATION()` functions, both of which return `binary` for a binary string argument:

```
mysql> SELECT CHARSET(BINARY 'a'), COLLATION(BINARY 'a');
+--------------------+----------------------+
| CHARSET(BINARY 'a') | COLLATION(BINARY 'a') |
```

```
+--------------------+----------------------+
| binary             | binary               |
+--------------------+----------------------+
```

For operations that combine multiple string inputs and return a single string output, the “aggregation rules” of standard SQL apply for determining the collation of the result:

- If an explicit `COLLATE Y` occurs, use *Y*.
- If explicit `COLLATE Y` and `COLLATE Z` occur, raise an error.
- Otherwise, if all collations are *Y*, use *Y*.
- Otherwise, the result has no collation.

For example, with `CASE ... WHEN a THEN b WHEN b THEN c COLLATE X END`, the resulting collation is *X*. The same applies for `UNION`, `||`, `CONCAT()`, `ELT()`, `GREATEST()`, `IF()`, and `LEAST()`.

For operations that convert to character data, the character set and collation of the strings that result from the operations are defined by the `character_set_connection` and `collation_connection` system variables that determine the default connection character set and collation (see Section 10.1.4, “Connection Character Sets and Collations”). This applies only to `BIN_TO_UUID()`, `CAST()`, `CONV()`, `FORMAT()`, `HEX()`, and `SPACE()`.

As of MySQL 8.0.2, an exception to the preceding priniciple occurs for expressions for virtual generated columns. In such expressions, the table character set is used for `BIN_TO_UUID()`, `CONV()`, or `HEX()` results, regardless of connection character set.

If there is any question about the character set or collation of the result returned by a string function, use the `CHARSET()` or `COLLATION()` function to find out:

```
mysql> SELECT USER(), CHARSET(USER()), COLLATION(USER());
+----------------+-----------------+-------------------+
| USER()         | CHARSET(USER()) | COLLATION(USER()) |
+----------------+-----------------+-------------------+
| test@localhost | utf8            | utf8_general_ci   |
+----------------+-----------------+-------------------+
mysql> SELECT CHARSET(COMPRESS('abc')), COLLATION(COMPRESS('abc'));
+--------------------------+----------------------------+
| CHARSET(COMPRESS('abc')) | COLLATION(COMPRESS('abc')) |
+--------------------------+----------------------------+
| binary                   | binary                     |
+--------------------------+----------------------------+
```

12.6 Numeric Functions and Operators

Table 12.10 Numeric Functions and Operators

Name	Description
`ABS()`	Return the absolute value
`ACOS()`	Return the arc cosine
`ASIN()`	Return the arc sine
`ATAN()`	Return the arc tangent
`ATAN2(), ATAN()`	Return the arc tangent of the two arguments
`CEIL()`	Return the smallest integer value not less than the argument
`CEILING()`	Return the smallest integer value not less than the argument
`CONV()`	Convert numbers between different number bases
`COS()`	Return the cosine

Name	Description
COT()	Return the cotangent
CRC32()	Compute a cyclic redundancy check value
DEGREES()	Convert radians to degrees
DIV	Integer division
/	Division operator
EXP()	Raise to the power of
FLOOR()	Return the largest integer value not greater than the argument
LN()	Return the natural logarithm of the argument
LOG()	Return the natural logarithm of the first argument
LOG10()	Return the base-10 logarithm of the argument
LOG2()	Return the base-2 logarithm of the argument
-	Minus operator
MOD()	Return the remainder
%, MOD	Modulo operator
PI()	Return the value of pi
+	Addition operator
POW()	Return the argument raised to the specified power
POWER()	Return the argument raised to the specified power
RADIANS()	Return argument converted to radians
RAND()	Return a random floating-point value
ROUND()	Round the argument
SIGN()	Return the sign of the argument
SIN()	Return the sine of the argument
SQRT()	Return the square root of the argument
TAN()	Return the tangent of the argument
*	Multiplication operator
TRUNCATE()	Truncate to specified number of decimal places
-	Change the sign of the argument

12.6.1 Arithmetic Operators

Table 12.11 Arithmetic Operators

Name	Description
DIV	Integer division
/	Division operator
-	Minus operator
%, MOD	Modulo operator
+	Addition operator
*	Multiplication operator
-	Change the sign of the argument

The usual arithmetic operators are available. The result is determined according to the following rules:

- In the case of `-`, `+`, and `*`, the result is calculated with `BIGINT` (64-bit) precision if both operands are integers.
- If both operands are integers and any of them are unsigned, the result is an unsigned integer. For subtraction, if the `NO_UNSIGNED_SUBTRACTION` SQL mode is enabled, the result is signed even if any operand is unsigned.
- If any of the operands of a `+`, `-`, `/`, `*`, `%` is a real or string value, the precision of the result is the precision of the operand with the maximum precision.
- In division performed with `/`, the scale of the result when using two exact-value operands is the scale of the first operand plus the value of the `div_precision_increment` system variable (which is 4 by default). For example, the result of the expression `5.05 / 0.014` has a scale of six decimal places (`360.714286`).

These rules are applied for each operation, such that nested calculations imply the precision of each component. Hence, `(14620 / 9432456) / (24250 / 9432456)`, resolves first to `(0.0014) / (0.0026)`, with the final result having 8 decimal places (`0.60288653`).

Because of these rules and the way they are applied, care should be taken to ensure that components and subcomponents of a calculation use the appropriate level of precision. See Section 12.10, “Cast Functions and Operators”.

For information about handling of overflow in numeric expression evaluation, see Section 11.2.6, “Out-of-Range and Overflow Handling”.

Arithmetic operators apply to numbers. For other types of values, alternative operations may be available. For example, to add date values, use `DATE_ADD()`; see Section 12.7, “Date and Time Functions”.

- `+`

  Addition:

```
mysql> SELECT 3+5;
        -> 8
```

- `-`

  Subtraction:

```
mysql> SELECT 3-5;
        -> -2
```

- `-`

  Unary minus. This operator changes the sign of the operand.

```
mysql> SELECT - 2;
        -> -2
```

Note

If this operator is used with a `BIGINT`, the return value is also a `BIGINT`. This means that you should avoid using `-` on integers that may have the value of -2^{63}.

- `*`

  Multiplication:

```
mysql> SELECT 3*5;
        -> 15
mysql> SELECT 18014398509481984*18014398509481984.0;
        -> 324518553658426726783156020576256.0
mysql> SELECT 18014398509481984*18014398509481984;
        -> out-of-range error
```

The last expression produces an error because the result of the integer multiplication exceeds the 64-bit range of `BIGINT` calculations. (See Section 11.2, “Numeric Types”.)

- `/`

  Division:

  ```
  mysql> SELECT 3/5;
          -> 0.60
  ```

  Division by zero produces a `NULL` result:

  ```
  mysql> SELECT 102/(1-1);
          -> NULL
  ```

  A division is calculated with `BIGINT` arithmetic only if performed in a context where its result is converted to an integer.

- `DIV`

  Integer division. Discards from the division result any fractional part to the right of the decimal point.

  If either operand has a noninteger type, the operands are converted to `DECIMAL` and divided using `DECIMAL` arithmetic before converting the result to `BIGINT`. If the result exceeds `BIGINT` range, an error occurs.

  ```
  mysql> SELECT 5 DIV 2, -5 DIV 2, 5 DIV -2, -5 DIV -2;
          -> 2, -2, -2, 2
  ```

- *`N`* `%` *`M`*, *`N`* `MOD` *`M`*

  Modulo operation. Returns the remainder of *`N`* divided by *`M`*. For more information, see the description for the `MOD()` function in Section 12.6.2, “Mathematical Functions”.

12.6.2 Mathematical Functions

Table 12.12 Mathematical Functions

Name	Description
`ABS()`	Return the absolute value
`ACOS()`	Return the arc cosine
`ASIN()`	Return the arc sine
`ATAN()`	Return the arc tangent
`ATAN2()`, `ATAN()`	Return the arc tangent of the two arguments
`CEIL()`	Return the smallest integer value not less than the argument
`CEILING()`	Return the smallest integer value not less than the argument
`CONV()`	Convert numbers between different number bases
`COS()`	Return the cosine
`COT()`	Return the cotangent
`CRC32()`	Compute a cyclic redundancy check value

Name	Description
DEGREES()	Convert radians to degrees
EXP()	Raise to the power of
FLOOR()	Return the largest integer value not greater than the argument
LN()	Return the natural logarithm of the argument
LOG()	Return the natural logarithm of the first argument
LOG10()	Return the base-10 logarithm of the argument
LOG2()	Return the base-2 logarithm of the argument
MOD()	Return the remainder
PI()	Return the value of pi
POW()	Return the argument raised to the specified power
POWER()	Return the argument raised to the specified power
RADIANS()	Return argument converted to radians
RAND()	Return a random floating-point value
ROUND()	Round the argument
SIGN()	Return the sign of the argument
SIN()	Return the sine of the argument
SQRT()	Return the square root of the argument
TAN()	Return the tangent of the argument
TRUNCATE()	Truncate to specified number of decimal places

All mathematical functions return NULL in the event of an error.

- ABS(*X*)

  Returns the absolute value of *X*.

```
mysql> SELECT ABS(2);
        -> 2
mysql> SELECT ABS(-32);
        -> 32
```

  This function is safe to use with BIGINT values.

- ACOS(*X*)

  Returns the arc cosine of *X*, that is, the value whose cosine is *X*. Returns NULL if *X* is not in the range -1 to 1.

```
mysql> SELECT ACOS(1);
        -> 0
mysql> SELECT ACOS(1.0001);
        -> NULL
mysql> SELECT ACOS(0);
        -> 1.5707963267949
```

- ASIN(*X*)

  Returns the arc sine of *X*, that is, the value whose sine is *X*. Returns NULL if *X* is not in the range -1 to 1.

```
mysql> SELECT ASIN(0.2);
```

```
       -> 0.20135792079033
mysql> SELECT ASIN('foo');

+-------------+
| ASIN('foo') |
+-------------+
|           0 |
+-------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS;
+---------+------+----------------------------------------+
| Level   | Code | Message                                |
+---------+------+----------------------------------------+
| Warning | 1292 | Truncated incorrect DOUBLE value: 'foo' |
+---------+------+----------------------------------------+
```

- `ATAN(X)`

  Returns the arc tangent of *X*, that is, the value whose tangent is *X*.

```
mysql> SELECT ATAN(2);
       -> 1.1071487177941
mysql> SELECT ATAN(-2);
       -> -1.1071487177941
```

- `ATAN(Y,X)`, `ATAN2(Y,X)`

  Returns the arc tangent of the two variables *X* and *Y*. It is similar to calculating the arc tangent of *Y* / *X*, except that the signs of both arguments are used to determine the quadrant of the result.

```
mysql> SELECT ATAN(-2,2);
       -> -0.78539816339745
mysql> SELECT ATAN2(PI(),0);
       -> 1.5707963267949
```

- `CEIL(X)`

  `CEIL()` is a synonym for `CEILING()`.

- `CEILING(X)`

  Returns the smallest integer value not less than *X*.

```
mysql> SELECT CEILING(1.23);
       -> 2
mysql> SELECT CEILING(-1.23);
       -> -1
```

  For exact-value numeric arguments, the return value has an exact-value numeric type. For string or floating-point arguments, the return value has a floating-point type.

- `CONV(N,from_base,to_base)`

  Converts numbers between different number bases. Returns a string representation of the number *N*, converted from base *from_base* to base *to_base*. Returns `NULL` if any argument is `NULL`. The argument *N* is interpreted as an integer, but may be specified as an integer or a string. The minimum base is `2` and the maximum base is `36`. If *from_base* is a negative number, *N* is regarded as a signed number. Otherwise, *N* is treated as unsigned. `CONV()` works with 64-bit precision.

```
mysql> SELECT CONV('a',16,2);
       -> '1010'
mysql> SELECT CONV('6E',18,8);
       -> '172'
```

```
mysql> SELECT CONV(-17,10,-18);
        -> '-H'
mysql> SELECT CONV(10+'10'+'10'+X'0a',10,10);
        -> '40'
```

- `COS(X)`

  Returns the cosine of *X*, where *X* is given in radians.

```
mysql> SELECT COS(PI());
        -> -1
```

- `COT(X)`

  Returns the cotangent of *X*.

```
mysql> SELECT COT(12);
        -> -1.5726734063977
mysql> SELECT COT(0);
        -> out-of-range error
```

- `CRC32(expr)`

  Computes a cyclic redundancy check value and returns a 32-bit unsigned value. The result is `NULL` if the argument is `NULL`. The argument is expected to be a string and (if possible) is treated as one if it is not.

```
mysql> SELECT CRC32('MySQL');
        -> 3259397556
mysql> SELECT CRC32('mysql');
        -> 2501908538
```

- `DEGREES(X)`

  Returns the argument *X*, converted from radians to degrees.

```
mysql> SELECT DEGREES(PI());
        -> 180
mysql> SELECT DEGREES(PI() / 2);
        -> 90
```

- `EXP(X)`

  Returns the value of *e* (the base of natural logarithms) raised to the power of *X*. The inverse of this function is `LOG()` (using a single argument only) or `LN()`.

```
mysql> SELECT EXP(2);
        -> 7.3890560989307
mysql> SELECT EXP(-2);
        -> 0.13533528323661
mysql> SELECT EXP(0);
        -> 1
```

- `FLOOR(X)`

  Returns the largest integer value not greater than *X*.

```
mysql> SELECT FLOOR(1.23), FLOOR(-1.23);
        -> 1, -2
```

  For exact-value numeric arguments, the return value has an exact-value numeric type. For string or floating-point arguments, the return value has a floating-point type.

- `FORMAT(X,D)`

  Formats the number *X* to a format like `'#,###,###.##'`, rounded to *D* decimal places, and returns the result as a string. For details, see Section 12.5, “String Functions”.

- `HEX(N_or_S)`

  This function can be used to obtain a hexadecimal representation of a decimal number or a string; the manner in which it does so varies according to the argument's type. See this function's description in Section 12.5, “String Functions”, for details.

- `LN(X)`

  Returns the natural logarithm of *X*; that is, the base-*e* logarithm of *X*. If *X* is less than or equal to 0.0E0, the function returns `NULL` and a warning “Invalid argument for logarithm” is reported.

```
mysql> SELECT LN(2);
        -> 0.69314718055995
mysql> SELECT LN(-2);
        -> NULL
```

  This function is synonymous with `LOG(X)`. The inverse of this function is the `EXP()` function.

- `LOG(X)`, `LOG(B,X)`

  If called with one parameter, this function returns the natural logarithm of *X*. If *X* is less than or equal to 0.0E0, the function returns `NULL` and a warning “Invalid argument for logarithm” is reported.

  The inverse of this function (when called with a single argument) is the `EXP()` function.

```
mysql> SELECT LOG(2);
        -> 0.69314718055995
mysql> SELECT LOG(-2);
        -> NULL
```

  If called with two parameters, this function returns the logarithm of *X* to the base *B*. If *X* is less than or equal to 0, or if *B* is less than or equal to 1, then `NULL` is returned.

```
mysql> SELECT LOG(2,65536);
        -> 16
mysql> SELECT LOG(10,100);
        -> 2
mysql> SELECT LOG(1,100);
        -> NULL
```

  `LOG(B,X)` is equivalent to `LOG(X) / LOG(B)`.

- `LOG2(X)`

  Returns the base-2 logarithm of *X*. If *X* is less than or equal to 0.0E0, the function returns `NULL` and a warning “Invalid argument for logarithm” is reported.

```
mysql> SELECT LOG2(65536);
        -> 16
mysql> SELECT LOG2(-100);
        -> NULL
```

  `LOG2()` is useful for finding out how many bits a number requires for storage. This function is equivalent to the expression `LOG(X) / LOG(2)`.

- `LOG10(X)`

Returns the base-10 logarithm of *X*. If *X* is less than or equal to 0.0E0, the function returns `NULL` and a warning “Invalid argument for logarithm” is reported.

```
mysql> SELECT LOG10(2);
        -> 0.30102999566398
mysql> SELECT LOG10(100);
        -> 2
mysql> SELECT LOG10(-100);
        -> NULL
```

`LOG10(X)` is equivalent to `LOG(10,X)`.

- `MOD(N,M)`, `N % M`, `N MOD M`

  Modulo operation. Returns the remainder of *N* divided by *M*.

  ```
  mysql> SELECT MOD(234, 10);
          -> 4
  mysql> SELECT 253 % 7;
          -> 1
  mysql> SELECT MOD(29,9);
          -> 2
  mysql> SELECT 29 MOD 9;
          -> 2
  ```

  This function is safe to use with `BIGINT` values.

  `MOD()` also works on values that have a fractional part and returns the exact remainder after division:

  ```
  mysql> SELECT MOD(34.5,3);
          -> 1.5
  ```

  `MOD(N,0)` returns `NULL`.

- `PI()`

  Returns the value of π (pi). The default number of decimal places displayed is seven, but MySQL uses the full double-precision value internally.

  ```
  mysql> SELECT PI();
          -> 3.141593
  mysql> SELECT PI()+0.00000000000000000;
          -> 3.141592653589793116
  ```

- `POW(X,Y)`

  Returns the value of *X* raised to the power of *Y*.

  ```
  mysql> SELECT POW(2,2);
          -> 4
  mysql> SELECT POW(2,-2);
          -> 0.25
  ```

- `POWER(X,Y)`

  This is a synonym for `POW()`.

- `RADIANS(X)`

  Returns the argument *X*, converted from degrees to radians. (Note that π radians equals 180 degrees.)

```
mysql> SELECT RADIANS(90);
       -> 1.5707963267949
```

- `RAND([N])`

  Returns a random floating-point value `v` in the range `0` <= `v` < `1.0`. To obtain a random integer `R` in the range `i` <= `R` < `j`, use the expression `FLOOR(i + RAND() * (j − i))`. For example, to obtain a random integer in the range the range `7` <= `R` < `12`, use the following statement:

```
SELECT FLOOR(7 + (RAND() * 5));
```

  If an integer argument `N` is specified, it is used as the seed value:

  - With a constant initializer argument, the seed is initialized once when the statement is prepared, prior to execution.
  - With a nonconstant initializer argument (such as a column name), the seed is initialized with the value for each invocation of `RAND()`.

  One implication of this behavior is that for equal argument values, `RAND(N)` returns the same value each time, and thus produces a repeatable sequence of column values. In the following example, the sequence of values produced by `RAND(3)` is the same both places it occurs.

```
mysql> CREATE TABLE t (i INT);
Query OK, 0 rows affected (0.42 sec)

mysql> INSERT INTO t VALUES(1),(2),(3);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT i, RAND() FROM t;
+------+------------------+
| i    | RAND()           |
+------+------------------+
|    1 | 0.61914388706828 |
|    2 | 0.93845168309142 |
|    3 | 0.83482678498591 |
+------+------------------+
3 rows in set (0.00 sec)

mysql> SELECT i, RAND(3) FROM t;
+------+------------------+
| i    | RAND(3)          |
+------+------------------+
|    1 | 0.90576975597606 |
|    2 | 0.37307905813035 |
|    3 | 0.14808605345719 |
+------+------------------+
3 rows in set (0.00 sec)

mysql> SELECT i, RAND() FROM t;
+------+------------------+
| i    | RAND()           |
+------+------------------+
|    1 | 0.35877890638893 |
|    2 | 0.28941420772058 |
|    3 | 0.37073435016976 |
+------+------------------+
3 rows in set (0.00 sec)

mysql> SELECT i, RAND(3) FROM t;
+------+------------------+
| i    | RAND(3)          |
+------+------------------+
```

```
|    1 | 0.90576975597606 |
|    2 | 0.37307905813035 |
|    3 | 0.14808605345719 |
+------+------------------+
3 rows in set (0.01 sec)
```

`RAND()` in a `WHERE` clause is evaluated for every row (when selecting from one table) or combination of rows (when selecting from a multiple-table join). Thus, for optimizer purposes, `RAND()` is not a constant value and cannot be used for index optimizations. For more information, see Section 8.2.1.17, “Function Call Optimization”.

Use of a column with `RAND()` values in an `ORDER BY` or `GROUP BY` clause may yield unexpected results because for either clause a `RAND()` expression can be evaluated multiple times for the same row, each time returning a different result. If the goal is to retrieve rows in random order, you can use a statement like this:

```
SELECT * FROM tbl_name ORDER BY RAND();
```

To select a random sample from a set of rows, combine `ORDER BY RAND()` with `LIMIT`:

```
SELECT * FROM table1, table2 WHERE a=b AND c<d ORDER BY RAND() LIMIT 1000;
```

`RAND()` is not meant to be a perfect random generator. It is a fast way to generate random numbers on demand that is portable between platforms for the same MySQL version.

This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`. (Bug #49222)

- `ROUND(X)`, `ROUND(X,D)`

  Rounds the argument *X* to *D* decimal places. The rounding algorithm depends on the data type of *X*. *D* defaults to 0 if not specified. *D* can be negative to cause *D* digits left of the decimal point of the value *X* to become zero.

```
mysql> SELECT ROUND(-1.23);
        -> -1
mysql> SELECT ROUND(-1.58);
        -> -2
mysql> SELECT ROUND(1.58);
        -> 2
mysql> SELECT ROUND(1.298, 1);
        -> 1.3
mysql> SELECT ROUND(1.298, 0);
        -> 1
mysql> SELECT ROUND(23.298, -1);
        -> 20
```

The return value has the same type as the first argument (assuming that it is integer, double, or decimal). This means that for an integer argument, the result is an integer (no decimal places):

```
mysql> SELECT ROUND(150.000,2), ROUND(150,2);
+------------------+--------------+
| ROUND(150.000,2) | ROUND(150,2) |
+------------------+--------------+
|           150.00 |          150 |
+------------------+--------------+
```

`ROUND()` uses the following rules depending on the type of the first argument:

- For exact-value numbers, `ROUND()` uses the “round half away from zero” or “round toward nearest” rule: A value with a fractional part of .5 or greater is rounded up to the next integer if positive or down to the next integer if negative. (In other words, it is rounded away from zero.) A

value with a fractional part less than .5 is rounded down to the next integer if positive or up to the next integer if negative.

- For approximate-value numbers, the result depends on the C library. On many systems, this means that `ROUND()` uses the "round to nearest even" rule: A value with any fractional part is rounded to the nearest even integer.

The following example shows how rounding differs for exact and approximate values:

```
mysql> SELECT ROUND(2.5), ROUND(25E-1);
+------------+--------------+
| ROUND(2.5) | ROUND(25E-1) |
+------------+--------------+
| 3          |            2 |
+------------+--------------+
```

For more information, see Section 12.21, “Precision Math”.

- `SIGN(X)`

  Returns the sign of the argument as `-1`, `0`, or `1`, depending on whether *X* is negative, zero, or positive.

```
mysql> SELECT SIGN(-32);
        -> -1
mysql> SELECT SIGN(0);
        -> 0
mysql> SELECT SIGN(234);
        -> 1
```

- `SIN(X)`

  Returns the sine of *X*, where *X* is given in radians.

```
mysql> SELECT SIN(PI());
        -> 1.2246063538224e-16
mysql> SELECT ROUND(SIN(PI()));
        -> 0
```

- `SQRT(X)`

  Returns the square root of a nonnegative number *X*.

```
mysql> SELECT SQRT(4);
        -> 2
mysql> SELECT SQRT(20);
        -> 4.4721359549996
mysql> SELECT SQRT(-16);
        -> NULL
```

- `TAN(X)`

  Returns the tangent of *X*, where *X* is given in radians.

```
mysql> SELECT TAN(PI());
        -> -1.2246063538224e-16
mysql> SELECT TAN(PI()+1);
        -> 1.5574077246549
```

- `TRUNCATE(X,D)`

Returns the number *X*, truncated to *D* decimal places. If *D* is `0`, the result has no decimal point or fractional part. *D* can be negative to cause *D* digits left of the decimal point of the value *X* to become zero.

```
mysql> SELECT TRUNCATE(1.223,1);
        -> 1.2
mysql> SELECT TRUNCATE(1.999,1);
        -> 1.9
mysql> SELECT TRUNCATE(1.999,0);
        -> 1
mysql> SELECT TRUNCATE(-1.999,1);
        -> -1.9
mysql> SELECT TRUNCATE(122,-2);
       -> 100
mysql> SELECT TRUNCATE(10.28*100,0);
       -> 1028
```

All numbers are rounded toward zero.

12.7 Date and Time Functions

This section describes the functions that can be used to manipulate temporal values. See Section 11.3, “Date and Time Types”, for a description of the range of values each date and time type has and the valid formats in which values may be specified.

Table 12.13 Date/Time Functions

Name	Description
`ADDDATE()`	Add time values (intervals) to a date value
`ADDTIME()`	Add time
`CONVERT_TZ()`	Convert from one time zone to another
`CURDATE()`	Return the current date
`CURRENT_DATE()`, `CURRENT_DATE`	Synonyms for CURDATE()
`CURRENT_TIME()`, `CURRENT_TIME`	Synonyms for CURTIME()
`CURRENT_TIMESTAMP()`, `CURRENT_TIMESTAMP`	Synonyms for NOW()
`CURTIME()`	Return the current time
`DATE()`	Extract the date part of a date or datetime expression
`DATE_ADD()`	Add time values (intervals) to a date value
`DATE_FORMAT()`	Format date as specified
`DATE_SUB()`	Subtract a time value (interval) from a date
`DATEDIFF()`	Subtract two dates
`DAY()`	Synonym for DAYOFMONTH()
`DAYNAME()`	Return the name of the weekday
`DAYOFMONTH()`	Return the day of the month (0-31)
`DAYOFWEEK()`	Return the weekday index of the argument
`DAYOFYEAR()`	Return the day of the year (1-366)
`EXTRACT()`	Extract part of a date
`FROM_DAYS()`	Convert a day number to a date
`FROM_UNIXTIME()`	Format Unix timestamp as a date
`GET_FORMAT()`	Return a date format string

Name	Description
HOUR()	Extract the hour
LAST_DAY	Return the last day of the month for the argument
LOCALTIME(), LOCALTIME	Synonym for NOW()
LOCALTIMESTAMP, LOCALTIMESTAMP()	Synonym for NOW()
MAKEDATE()	Create a date from the year and day of year
MAKETIME()	Create time from hour, minute, second
MICROSECOND()	Return the microseconds from argument
MINUTE()	Return the minute from the argument
MONTH()	Return the month from the date passed
MONTHNAME()	Return the name of the month
NOW()	Return the current date and time
PERIOD_ADD()	Add a period to a year-month
PERIOD_DIFF()	Return the number of months between periods
QUARTER()	Return the quarter from a date argument
SEC_TO_TIME()	Converts seconds to 'HH:MM:SS' format
SECOND()	Return the second (0-59)
STR_TO_DATE()	Convert a string to a date
SUBDATE()	Synonym for DATE_SUB() when invoked with three arguments
SUBTIME()	Subtract times
SYSDATE()	Return the time at which the function executes
TIME()	Extract the time portion of the expression passed
TIME_FORMAT()	Format as time
TIME_TO_SEC()	Return the argument converted to seconds
TIMEDIFF()	Subtract time
TIMESTAMP()	With a single argument, this function returns the date or datetime expression; with two arguments, the sum of the arguments
TIMESTAMPADD()	Add an interval to a datetime expression
TIMESTAMPDIFF()	Subtract an interval from a datetime expression
TO_DAYS()	Return the date argument converted to days
TO_SECONDS()	Return the date or datetime argument converted to seconds since Year 0
UNIX_TIMESTAMP()	Return a Unix timestamp
UTC_DATE()	Return the current UTC date
UTC_TIME()	Return the current UTC time
UTC_TIMESTAMP()	Return the current UTC date and time
WEEK()	Return the week number
WEEKDAY()	Return the weekday index
WEEKOFYEAR()	Return the calendar week of the date (1-53)
YEAR()	Return the year

Name	Description
YEARWEEK()	Return the year and week

Here is an example that uses date functions. The following query selects all rows with a *date_col* value from within the last 30 days:

```
mysql> SELECT something FROM tbl_name
    -> WHERE DATE_SUB(CURDATE(),INTERVAL 30 DAY) <= date_col;
```

The query also selects rows with dates that lie in the future.

Functions that expect date values usually accept datetime values and ignore the time part. Functions that expect time values usually accept datetime values and ignore the date part.

Functions that return the current date or time each are evaluated only once per query at the start of query execution. This means that multiple references to a function such as NOW() within a single query always produce the same result. (For our purposes, a single query also includes a call to a stored program (stored routine, trigger, or event) and all subprograms called by that program.) This principle also applies to CURDATE(), CURTIME(), UTC_DATE(), UTC_TIME(), UTC_TIMESTAMP(), and to any of their synonyms.

The CURRENT_TIMESTAMP(), CURRENT_TIME(), CURRENT_DATE(), and FROM_UNIXTIME() functions return values in the connection's current time zone, which is available as the value of the time_zone system variable. In addition, UNIX_TIMESTAMP() assumes that its argument is a datetime value in the current time zone. See Section 10.6, “MySQL Server Time Zone Support”.

Some date functions can be used with “zero” dates or incomplete dates such as '2001-11-00', whereas others cannot. Functions that extract parts of dates typically work with incomplete dates and thus can return 0 when you might otherwise expect a nonzero value. For example:

```
mysql> SELECT DAYOFMONTH('2001-11-00'), MONTH('2005-00-00');
        -> 0, 0
```

Other functions expect complete dates and return NULL for incomplete dates. These include functions that perform date arithmetic or that map parts of dates to names. For example:

```
mysql> SELECT DATE_ADD('2006-05-00',INTERVAL 1 DAY);
        -> NULL
mysql> SELECT DAYNAME('2006-05-00');
        -> NULL
```

Several functions are more strict when passed a DATE() function value as their argument and reject incomplete dates with a day part of zero. These functions are affected: CONVERT_TZ(), DATE_ADD(), DATE_SUB(), DAYOFYEAR(), LAST_DAY() (permits a day part of zero), TIMESTAMPDIFF(), TO_DAYS(), TO_SECONDS(), WEEK(), WEEKDAY(), WEEKOFYEAR(), YEARWEEK().

Fractional seconds for TIME, DATETIME, and TIMESTAMP values are supported, with up to microsecond precision. Functions that take temporal arguments accept values with fractional seconds. Return values from temporal functions include fractional seconds as appropriate.

- ADDDATE(*date*,INTERVAL *expr unit*), ADDDATE(*expr*,*days*)

  When invoked with the INTERVAL form of the second argument, ADDDATE() is a synonym for DATE_ADD(). The related function SUBDATE() is a synonym for DATE_SUB(). For information on the INTERVAL *unit* argument, see the discussion for DATE_ADD().

  ```
  mysql> SELECT DATE_ADD('2008-01-02', INTERVAL 31 DAY);
          -> '2008-02-02'
  mysql> SELECT ADDDATE('2008-01-02', INTERVAL 31 DAY);
          -> '2008-02-02'
  ```

When invoked with the *days* form of the second argument, MySQL treats it as an integer number of days to be added to *expr*.

```
mysql> SELECT ADDDATE('2008-01-02', 31);
        -> '2008-02-02'
```

- `ADDTIME(expr1,expr2)`

  `ADDTIME()` adds *expr2* to *expr1* and returns the result. *expr1* is a time or datetime expression, and *expr2* is a time expression.

```
mysql> SELECT ADDTIME('2007-12-31 23:59:59.999999', '1 1:1:1.000002');
        -> '2008-01-02 01:01:01.000001'
mysql> SELECT ADDTIME('01:00:00.999999', '02:00:00.999998');
        -> '03:00:01.999997'
```

- `CONVERT_TZ(dt,from_tz,to_tz)`

  `CONVERT_TZ()` converts a datetime value *dt* from the time zone given by *from_tz* to the time zone given by *to_tz* and returns the resulting value. Time zones are specified as described in Section 10.6, “MySQL Server Time Zone Support”. This function returns `NULL` if the arguments are invalid.

  If the value falls out of the supported range of the `TIMESTAMP` type when converted from *from_tz* to UTC, no conversion occurs. The `TIMESTAMP` range is described in Section 11.1.2, “Date and Time Type Overview”.

```
mysql> SELECT CONVERT_TZ('2004-01-01 12:00:00','GMT','MET');
        -> '2004-01-01 13:00:00'
mysql> SELECT CONVERT_TZ('2004-01-01 12:00:00','+00:00','+10:00');
        -> '2004-01-01 22:00:00'
```

Note

To use named time zones such as `'MET'` or `'Europe/Moscow'`, the time zone tables must be properly set up. See Section 10.6, “MySQL Server Time Zone Support”, for instructions.

- `CURDATE()`

  Returns the current date as a value in `'YYYY-MM-DD'` or `YYYYMMDD` format, depending on whether the function is used in a string or numeric context.

```
mysql> SELECT CURDATE();
        -> '2008-06-13'
mysql> SELECT CURDATE() + 0;
        -> 20080613
```

- `CURRENT_DATE`, `CURRENT_DATE()`

  `CURRENT_DATE` and `CURRENT_DATE()` are synonyms for `CURDATE()`.

- `CURRENT_TIME`, `CURRENT_TIME([fsp])`

  `CURRENT_TIME` and `CURRENT_TIME()` are synonyms for `CURTIME()`.

- `CURRENT_TIMESTAMP`, `CURRENT_TIMESTAMP([fsp])`

  `CURRENT_TIMESTAMP` and `CURRENT_TIMESTAMP()` are synonyms for `NOW()`.

- `CURTIME([fsp])`

Returns the current time as a value in `'HH:MM:SS'` or `HHMMSS` format, depending on whether the function is used in a string or numeric context. The value is expressed in the current time zone.

If the *`fsp`* argument is given to specify a fractional seconds precision from 0 to 6, the return value includes a fractional seconds part of that many digits.

```
mysql> SELECT CURTIME();
        -> '23:50:26'
mysql> SELECT CURTIME() + 0;
        -> 235026.000000
```

- `DATE(`*`expr`*`)`

  Extracts the date part of the date or datetime expression *`expr`*.

  ```
  mysql> SELECT DATE('2003-12-31 01:02:03');
          -> '2003-12-31'
  ```

- `DATEDIFF(`*`expr1`*`,`*`expr2`*`)`

  `DATEDIFF()` returns *`expr1`* − *`expr2`* expressed as a value in days from one date to the other. *`expr1`* and *`expr2`* are date or date-and-time expressions. Only the date parts of the values are used in the calculation.

  ```
  mysql> SELECT DATEDIFF('2007-12-31 23:59:59','2007-12-30');
          -> 1
  mysql> SELECT DATEDIFF('2010-11-30 23:59:59','2010-12-31');
          -> -31
  ```

- `DATE_ADD(`*`date`*`,INTERVAL `*`expr unit`*`)`, `DATE_SUB(`*`date`*`,INTERVAL `*`expr unit`*`)`

  These functions perform date arithmetic. The *`date`* argument specifies the starting date or datetime value. *`expr`* is an expression specifying the interval value to be added or subtracted from the starting date. *`expr`* is a string; it may start with a `-` for negative intervals. *`unit`* is a keyword indicating the units in which the expression should be interpreted.

  The `INTERVAL` keyword and the *`unit`* specifier are not case sensitive.

  The following table shows the expected form of the *`expr`* argument for each *`unit`* value.

***`unit`* Value**	**Expected *`expr`* Format**
`MICROSECOND`	`MICROSECONDS`
`SECOND`	`SECONDS`
`MINUTE`	`MINUTES`
`HOUR`	`HOURS`
`DAY`	`DAYS`
`WEEK`	`WEEKS`
`MONTH`	`MONTHS`
`QUARTER`	`QUARTERS`
`YEAR`	`YEARS`
`SECOND_MICROSECOND`	`'SECONDS.MICROSECONDS'`
`MINUTE_MICROSECOND`	`'MINUTES:SECONDS.MICROSECONDS'`
`MINUTE_SECOND`	`'MINUTES:SECONDS'`
`HOUR_MICROSECOND`	`'HOURS:MINUTES:SECONDS.MICROSECONDS'`

unit Value	Expected *expr* Format
HOUR_SECOND	'HOURS:MINUTES:SECONDS'
HOUR_MINUTE	'HOURS:MINUTES'
DAY_MICROSECOND	'DAYS HOURS:MINUTES:SECONDS.MICROSECONDS'
DAY_SECOND	'DAYS HOURS:MINUTES:SECONDS'
DAY_MINUTE	'DAYS HOURS:MINUTES'
DAY_HOUR	'DAYS HOURS'
YEAR_MONTH	'YEARS-MONTHS'

The return value depends on the arguments:

- DATETIME if the first argument is a DATETIME (or TIMESTAMP) value, or if the first argument is a DATE and the *unit* value uses HOURS, MINUTES, or SECONDS.
- String otherwise.

To ensure that the result is DATETIME, you can use CAST() to convert the first argument to DATETIME.

MySQL permits any punctuation delimiter in the *expr* format. Those shown in the table are the suggested delimiters. If the *date* argument is a DATE value and your calculations involve only YEAR, MONTH, and DAY parts (that is, no time parts), the result is a DATE value. Otherwise, the result is a DATETIME value.

Date arithmetic also can be performed using INTERVAL together with the + or - operator:

```
date + INTERVAL expr unit
date - INTERVAL expr unit
```

INTERVAL *expr unit* is permitted on either side of the + operator if the expression on the other side is a date or datetime value. For the - operator, INTERVAL *expr unit* is permitted only on the right side, because it makes no sense to subtract a date or datetime value from an interval.

```
mysql> SELECT '2008-12-31 23:59:59' + INTERVAL 1 SECOND;
        -> '2009-01-01 00:00:00'
mysql> SELECT INTERVAL 1 DAY + '2008-12-31';
        -> '2009-01-01'
mysql> SELECT '2005-01-01' - INTERVAL 1 SECOND;
        -> '2004-12-31 23:59:59'
mysql> SELECT DATE_ADD('2000-12-31 23:59:59',
    ->                 INTERVAL 1 SECOND);
        -> '2001-01-01 00:00:00'
mysql> SELECT DATE_ADD('2010-12-31 23:59:59',
    ->                 INTERVAL 1 DAY);
        -> '2011-01-01 23:59:59'
mysql> SELECT DATE_ADD('2100-12-31 23:59:59',
    ->                 INTERVAL '1:1' MINUTE_SECOND);
        -> '2101-01-01 00:01:00'
mysql> SELECT DATE_SUB('2005-01-01 00:00:00',
    ->                 INTERVAL '1 1:1:1' DAY_SECOND);
        -> '2004-12-30 22:58:59'
mysql> SELECT DATE_ADD('1900-01-01 00:00:00',
    ->                 INTERVAL '-1 10' DAY_HOUR);
        -> '1899-12-30 14:00:00'
mysql> SELECT DATE_SUB('1998-01-02', INTERVAL 31 DAY);
        -> '1997-12-02'
mysql> SELECT DATE_ADD('1992-12-31 23:59:59.000002',
    ->            INTERVAL '1.999999' SECOND_MICROSECOND);
        -> '1993-01-01 00:00:01.000001'
```

If you specify an interval value that is too short (does not include all the interval parts that would be expected from the *unit* keyword), MySQL assumes that you have left out the leftmost parts of the interval value. For example, if you specify a *unit* of DAY_SECOND, the value of *expr* is expected to have days, hours, minutes, and seconds parts. If you specify a value like '1:10', MySQL assumes that the days and hours parts are missing and the value represents minutes and seconds. In other words, '1:10' DAY_SECOND is interpreted in such a way that it is equivalent to '1:10' MINUTE_SECOND. This is analogous to the way that MySQL interprets TIME values as representing elapsed time rather than as a time of day.

Because *expr* is treated as a string, be careful if you specify a nonstring value with INTERVAL. For example, with an interval specifier of HOUR_MINUTE, 6/4 evaluates to 1.5000 and is treated as 1 hour, 5000 minutes:

```
mysql> SELECT 6/4;
        -> 1.5000
mysql> SELECT DATE_ADD('2009-01-01', INTERVAL 6/4 HOUR_MINUTE);
        -> '2009-01-04 12:20:00'
```

To ensure interpretation of the interval value as you expect, a CAST() operation may be used. To treat 6/4 as 1 hour, 5 minutes, cast it to a DECIMAL value with a single fractional digit:

```
mysql> SELECT CAST(6/4 AS DECIMAL(3,1));
        -> 1.5
mysql> SELECT DATE_ADD('1970-01-01 12:00:00',
    ->                 INTERVAL CAST(6/4 AS DECIMAL(3,1)) HOUR_MINUTE);
        -> '1970-01-01 13:05:00'
```

If you add to or subtract from a date value something that contains a time part, the result is automatically converted to a datetime value:

```
mysql> SELECT DATE_ADD('2013-01-01', INTERVAL 1 DAY);
        -> '2013-01-02'
mysql> SELECT DATE_ADD('2013-01-01', INTERVAL 1 HOUR);
        -> '2013-01-01 01:00:00'
```

If you add MONTH, YEAR_MONTH, or YEAR and the resulting date has a day that is larger than the maximum day for the new month, the day is adjusted to the maximum days in the new month:

```
mysql> SELECT DATE_ADD('2009-01-30', INTERVAL 1 MONTH);
        -> '2009-02-28'
```

Date arithmetic operations require complete dates and do not work with incomplete dates such as '2006-07-00' or badly malformed dates:

```
mysql> SELECT DATE_ADD('2006-07-00', INTERVAL 1 DAY);
        -> NULL
mysql> SELECT '2005-03-32' + INTERVAL 1 MONTH;
        -> NULL
```

- DATE_FORMAT(*date*,*format*)

  Formats the *date* value according to the *format* string.

  The following specifiers may be used in the *format* string. The % character is required before format specifier characters.

Specifier	Description
%a	Abbreviated weekday name (Sun..Sat)

Specifier	Description
%b	Abbreviated month name (Jan..Dec)
%c	Month, numeric (0..12)
%D	Day of the month with English suffix (0th, 1st, 2nd, 3rd, …)
%d	Day of the month, numeric (00..31)
%e	Day of the month, numeric (0..31)
%f	Microseconds (000000..999999)
%H	Hour (00..23)
%h	Hour (01..12)
%I	Hour (01..12)
%i	Minutes, numeric (00..59)
%j	Day of year (001..366)
%k	Hour (0..23)
%l	Hour (1..12)
%M	Month name (January..December)
%m	Month, numeric (00..12)
%p	AM or PM
%r	Time, 12-hour (hh:mm:ss followed by AM or PM)
%S	Seconds (00..59)
%s	Seconds (00..59)
%T	Time, 24-hour (hh:mm:ss)
%U	Week (00..53), where Sunday is the first day of the week; WEEK() mode 0
%u	Week (00..53), where Monday is the first day of the week; WEEK() mode 1
%V	Week (01..53), where Sunday is the first day of the week; WEEK() mode 2; used with %X
%v	Week (01..53), where Monday is the first day of the week; WEEK() mode 3; used with %x
%W	Weekday name (Sunday..Saturday)
%w	Day of the week (0=Sunday..6=Saturday)
%X	Year for the week where Sunday is the first day of the week, numeric, four digits; used with %V
%x	Year for the week, where Monday is the first day of the week, numeric, four digits; used with %v
%Y	Year, numeric, four digits
%y	Year, numeric (two digits)
%%	A literal % character
%*x*	*x*, for any "*x*" not listed above

Ranges for the month and day specifiers begin with zero due to the fact that MySQL permits the storing of incomplete dates such as '2014-00-00'.

The language used for day and month names and abbreviations is controlled by the value of the lc_time_names system variable (Section 10.7, "MySQL Server Locale Support").

For the %U, %u, %V, and %v specifiers, see the description of the WEEK() function for information about the mode values. The mode affects how week numbering occurs.

`DATE_FORMAT()` returns a string with a character set and collation given by `character_set_connection` and `collation_connection` so that it can return month and weekday names containing non-ASCII characters.

```
mysql> SELECT DATE_FORMAT('2009-10-04 22:23:00', '%W %M %Y');
        -> 'Sunday October 2009'
mysql> SELECT DATE_FORMAT('2007-10-04 22:23:00', '%H:%i:%s');
        -> '22:23:00'
mysql> SELECT DATE_FORMAT('1900-10-04 22:23:00',
    ->                 '%D %y %a %d %m %b %j');
        -> '4th 00 Thu 04 10 Oct 277'
mysql> SELECT DATE_FORMAT('1997-10-04 22:23:00',
    ->                 '%H %k %I %r %T %S %w');
        -> '22 22 10 10:23:00 PM 22:23:00 00 6'
mysql> SELECT DATE_FORMAT('1999-01-01', '%X %V');
        -> '1998 52'
mysql> SELECT DATE_FORMAT('2006-06-00', '%d');
        -> '00'
```

- `DATE_SUB(date,INTERVAL expr unit)`

  See the description for `DATE_ADD()`.

- `DAY(date)`

  `DAY()` is a synonym for `DAYOFMONTH()`.

- `DAYNAME(date)`

  Returns the name of the weekday for *`date`*. The language used for the name is controlled by the value of the `lc_time_names` system variable (Section 10.7, “MySQL Server Locale Support”).

  ```
  mysql> SELECT DAYNAME('2007-02-03');
          -> 'Saturday'
  ```

- `DAYOFMONTH(date)`

  Returns the day of the month for *`date`*, in the range `1` to `31`, or `0` for dates such as `'0000-00-00'` or `'2008-00-00'` that have a zero day part.

  ```
  mysql> SELECT DAYOFMONTH('2007-02-03');
          -> 3
  ```

- `DAYOFWEEK(date)`

  Returns the weekday index for *`date`* (`1` = Sunday, `2` = Monday, …, `7` = Saturday). These index values correspond to the ODBC standard.

  ```
  mysql> SELECT DAYOFWEEK('2007-02-03');
          -> 7
  ```

- `DAYOFYEAR(date)`

  Returns the day of the year for *`date`*, in the range `1` to `366`.

  ```
  mysql> SELECT DAYOFYEAR('2007-02-03');
          -> 34
  ```

- `EXTRACT(unit FROM date)`

  The `EXTRACT()` function uses the same kinds of unit specifiers as `DATE_ADD()` or `DATE_SUB()`, but extracts parts from the date rather than performing date arithmetic.

```
mysql> SELECT EXTRACT(YEAR FROM '2009-07-02');
       -> 2009
mysql> SELECT EXTRACT(YEAR_MONTH FROM '2009-07-02 01:02:03');
       -> 200907
mysql> SELECT EXTRACT(DAY_MINUTE FROM '2009-07-02 01:02:03');
       -> 20102
mysql> SELECT EXTRACT(MICROSECOND
    ->                FROM '2003-01-02 10:30:00.000123');
        -> 123
```

- `FROM_DAYS(N)`

  Given a day number *`N`*, returns a `DATE` value.

```
mysql> SELECT FROM_DAYS(730669);
        -> '2007-07-03'
```

  Use `FROM_DAYS()` with caution on old dates. It is not intended for use with values that precede the advent of the Gregorian calendar (1582). See Section 12.8, “What Calendar Is Used By MySQL?”.

- `FROM_UNIXTIME(unix_timestamp)`, `FROM_UNIXTIME(unix_timestamp,format)`

  Returns a representation of the *`unix_timestamp`* argument as a value in `'YYYY-MM-DD HH:MM:SS'` or `YYYYMMDDHHMMSS` format, depending on whether the function is used in a string or numeric context. The value is expressed in the current time zone. *`unix_timestamp`* is an internal timestamp value such as is produced by the `UNIX_TIMESTAMP()` function.

  If *`format`* is given, the result is formatted according to the *`format`* string, which is used the same way as listed in the entry for the `DATE_FORMAT()` function.

```
mysql> SELECT FROM_UNIXTIME(1447430881);
        -> '2015-11-13 10:08:01'
mysql> SELECT FROM_UNIXTIME(1447430881) + 0;
        -> 20151113100801
mysql> SELECT FROM_UNIXTIME(UNIX_TIMESTAMP(),
    ->                      '%Y %D %M %h:%i:%s %x');
        -> '2015 13th November 10:08:01 2015'
```

  Note: If you use `UNIX_TIMESTAMP()` and `FROM_UNIXTIME()` to convert between `TIMESTAMP` values and Unix timestamp values, the conversion is lossy because the mapping is not one-to-one in both directions. For details, see the description of the `UNIX_TIMESTAMP()` function.

- `GET_FORMAT({DATE|TIME|DATETIME}, {'EUR'|'USA'|'JIS'|'ISO'|'INTERNAL'})`

  Returns a format string. This function is useful in combination with the `DATE_FORMAT()` and the `STR_TO_DATE()` functions.

  The possible values for the first and second arguments result in several possible format strings (for the specifiers used, see the table in the `DATE_FORMAT()` function description). ISO format refers to ISO 9075, not ISO 8601.

Function Call	Result
`GET_FORMAT(DATE,'USA')`	`'%m.%d.%Y'`
`GET_FORMAT(DATE,'JIS')`	`'%Y-%m-%d'`
`GET_FORMAT(DATE,'ISO')`	`'%Y-%m-%d'`
`GET_FORMAT(DATE,'EUR')`	`'%d.%m.%Y'`
`GET_FORMAT(DATE,'INTERNAL')`	`'%Y%m%d'`
`GET_FORMAT(DATETIME,'USA')`	`'%Y-%m-%d %H.%i.%s'`

Function Call	Result
GET_FORMAT(DATETIME,'JIS')	'%Y-%m-%d %H:%i:%s'
GET_FORMAT(DATETIME,'ISO')	'%Y-%m-%d %H:%i:%s'
GET_FORMAT(DATETIME,'EUR')	'%Y-%m-%d %H.%i.%s'
GET_FORMAT(DATETIME,'INTERNAL')	'%Y%m%d%H%i%s'
GET_FORMAT(TIME,'USA')	'%h:%i:%s %p'
GET_FORMAT(TIME,'JIS')	'%H:%i:%s'
GET_FORMAT(TIME,'ISO')	'%H:%i:%s'
GET_FORMAT(TIME,'EUR')	'%H.%i.%s'
GET_FORMAT(TIME,'INTERNAL')	'%H%i%s'

TIMESTAMP can also be used as the first argument to GET_FORMAT(), in which case the function returns the same values as for DATETIME.

```
mysql> SELECT DATE_FORMAT('2003-10-03',GET_FORMAT(DATE,'EUR'));
        -> '03.10.2003'
mysql> SELECT STR_TO_DATE('10.31.2003',GET_FORMAT(DATE,'USA'));
        -> '2003-10-31'
```

- HOUR(*time*)

  Returns the hour for *time*. The range of the return value is 0 to 23 for time-of-day values. However, the range of TIME values actually is much larger, so HOUR can return values greater than 23.

```
mysql> SELECT HOUR('10:05:03');
        -> 10
mysql> SELECT HOUR('272:59:59');
        -> 272
```

- LAST_DAY(*date*)

  Takes a date or datetime value and returns the corresponding value for the last day of the month. Returns NULL if the argument is invalid.

```
mysql> SELECT LAST_DAY('2003-02-05');
        -> '2003-02-28'
mysql> SELECT LAST_DAY('2004-02-05');
        -> '2004-02-29'
mysql> SELECT LAST_DAY('2004-01-01 01:01:01');
        -> '2004-01-31'
mysql> SELECT LAST_DAY('2003-03-32');
        -> NULL
```

- LOCALTIME, LOCALTIME([*fsp*])

  LOCALTIME and LOCALTIME() are synonyms for NOW().

- LOCALTIMESTAMP, LOCALTIMESTAMP([*fsp*])

  LOCALTIMESTAMP and LOCALTIMESTAMP() are synonyms for NOW().

- MAKEDATE(*year*,*dayofyear*)

  Returns a date, given year and day-of-year values. *dayofyear* must be greater than 0 or the result is NULL.

```
mysql> SELECT MAKEDATE(2011,31), MAKEDATE(2011,32);
        -> '2011-01-31', '2011-02-01'
```

```
mysql> SELECT MAKEDATE(2011,365), MAKEDATE(2014,365);
        -> '2011-12-31', '2014-12-31'
mysql> SELECT MAKEDATE(2011,0);
        -> NULL
```

- `MAKETIME(hour,minute,second)`

  Returns a time value calculated from the *`hour`*, *`minute`*, and *`second`* arguments.

  The *`second`* argument can have a fractional part.

```
mysql> SELECT MAKETIME(12,15,30);
        -> '12:15:30'
```

- `MICROSECOND(expr)`

  Returns the microseconds from the time or datetime expression *`expr`* as a number in the range from `0` to `999999`.

```
mysql> SELECT MICROSECOND('12:00:00.123456');
        -> 123456
mysql> SELECT MICROSECOND('2009-12-31 23:59:59.000010');
        -> 10
```

- `MINUTE(time)`

  Returns the minute for *`time`*, in the range `0` to `59`.

```
mysql> SELECT MINUTE('2008-02-03 10:05:03');
        -> 5
```

- `MONTH(date)`

  Returns the month for *`date`*, in the range `1` to `12` for January to December, or `0` for dates such as `'0000-00-00'` or `'2008-00-00'` that have a zero month part.

```
mysql> SELECT MONTH('2008-02-03');
        -> 2
```

- `MONTHNAME(date)`

  Returns the full name of the month for *`date`*. The language used for the name is controlled by the value of the `lc_time_names` system variable (Section 10.7, “MySQL Server Locale Support”).

```
mysql> SELECT MONTHNAME('2008-02-03');
        -> 'February'
```

- `NOW([fsp])`

  Returns the current date and time as a value in `'YYYY-MM-DD HH:MM:SS'` or `YYYYMMDDHHMMSS` format, depending on whether the function is used in a string or numeric context. The value is expressed in the current time zone.

  If the *`fsp`* argument is given to specify a fractional seconds precision from 0 to 6, the return value includes a fractional seconds part of that many digits.

```
mysql> SELECT NOW();
        -> '2007-12-15 23:50:26'
mysql> SELECT NOW() + 0;
        -> 20071215235026.000000
```

`NOW()` returns a constant time that indicates the time at which the statement began to execute. (Within a stored function or trigger, `NOW()` returns the time at which the function or triggering statement began to execute.) This differs from the behavior for `SYSDATE()`, which returns the exact time at which it executes.

```
mysql> SELECT NOW(), SLEEP(2), NOW();
+---------------------+----------+---------------------+
| NOW()               | SLEEP(2) | NOW()               |
+---------------------+----------+---------------------+
| 2006-04-12 13:47:36 |        0 | 2006-04-12 13:47:36 |
+---------------------+----------+---------------------+

mysql> SELECT SYSDATE(), SLEEP(2), SYSDATE();
+---------------------+----------+---------------------+
| SYSDATE()           | SLEEP(2) | SYSDATE()           |
+---------------------+----------+---------------------+
| 2006-04-12 13:47:44 |        0 | 2006-04-12 13:47:46 |
+---------------------+----------+---------------------+
```

In addition, the `SET TIMESTAMP` statement affects the value returned by `NOW()` but not by `SYSDATE()`. This means that timestamp settings in the binary log have no effect on invocations of `SYSDATE()`. Setting the timestamp to a nonzero value causes each subsequent invocation of `NOW()` to return that value. Setting the timestamp to zero cancels this effect so that `NOW()` once again returns the current date and time.

See the description for `SYSDATE()` for additional information about the differences between the two functions.

- `PERIOD_ADD(P,N)`

  Adds *`N`* months to period *`P`* (in the format `YYMM` or `YYYYMM`). Returns a value in the format `YYYYMM`. Note that the period argument *`P`* is *not* a date value.

```
mysql> SELECT PERIOD_ADD(200801,2);
        -> 200803
```

- `PERIOD_DIFF(P1,P2)`

  Returns the number of months between periods *`P1`* and *`P2`*. *`P1`* and *`P2`* should be in the format `YYMM` or `YYYYMM`. Note that the period arguments *`P1`* and *`P2`* are *not* date values.

```
mysql> SELECT PERIOD_DIFF(200802,200703);
        -> 11
```

- `QUARTER(date)`

  Returns the quarter of the year for *`date`*, in the range `1` to `4`.

```
mysql> SELECT QUARTER('2008-04-01');
        -> 2
```

- `SECOND(time)`

  Returns the second for *`time`*, in the range `0` to `59`.

```
mysql> SELECT SECOND('10:05:03');
        -> 3
```

- `SEC_TO_TIME(seconds)`

Returns the *`seconds`* argument, converted to hours, minutes, and seconds, as a `TIME` value. The range of the result is constrained to that of the `TIME` data type. A warning occurs if the argument corresponds to a value outside that range.

```
mysql> SELECT SEC_TO_TIME(2378);
        -> '00:39:38'
mysql> SELECT SEC_TO_TIME(2378) + 0;
        -> 3938
```

- `STR_TO_DATE(`*`str`*`,`*`format`*`)`

This is the inverse of the `DATE_FORMAT()` function. It takes a string *`str`* and a format string *`format`*. `STR_TO_DATE()` returns a `DATETIME` value if the format string contains both date and time parts, or a `DATE` or `TIME` value if the string contains only date or time parts. If the date, time, or datetime value extracted from *`str`* is illegal, `STR_TO_DATE()` returns `NULL` and produces a warning.

The server scans *`str`* attempting to match *`format`* to it. The format string can contain literal characters and format specifiers beginning with `%`. Literal characters in *`format`* must match literally in *`str`*. Format specifiers in *`format`* must match a date or time part in *`str`*. For the specifiers that can be used in *`format`*, see the `DATE_FORMAT()` function description.

```
mysql> SELECT STR_TO_DATE('01,5,2013','%d,%m,%Y');
        -> '2013-05-01'
mysql> SELECT STR_TO_DATE('May 1, 2013','%M %d,%Y');
        -> '2013-05-01'
```

Scanning starts at the beginning of *`str`* and fails if *`format`* is found not to match. Extra characters at the end of *`str`* are ignored.

```
mysql> SELECT STR_TO_DATE('a09:30:17','a%h:%i:%s');
        -> '09:30:17'
mysql> SELECT STR_TO_DATE('a09:30:17','%h:%i:%s');
        -> NULL
mysql> SELECT STR_TO_DATE('09:30:17a','%h:%i:%s');
        -> '09:30:17'
```

Unspecified date or time parts have a value of 0, so incompletely specified values in *`str`* produce a result with some or all parts set to 0:

```
mysql> SELECT STR_TO_DATE('abc','abc');
        -> '0000-00-00'
mysql> SELECT STR_TO_DATE('9','%m');
        -> '0000-09-00'
mysql> SELECT STR_TO_DATE('9','%s');
        -> '00:00:09'
```

Range checking on the parts of date values is as described in Section 11.3.1, “The DATE, DATETIME, and TIMESTAMP Types”. This means, for example, that “zero” dates or dates with part values of 0 are permitted unless the SQL mode is set to disallow such values.

```
mysql> SELECT STR_TO_DATE('00/00/0000', '%m/%d/%Y');
        -> '0000-00-00'
mysql> SELECT STR_TO_DATE('04/31/2004', '%m/%d/%Y');
        -> '2004-04-31'
```

If the `NO_ZERO_DATE` or `NO_ZERO_IN_DATE` SQL mode is enabled, zero dates or part of dates are disallowed. In that case, `STR_TO_DATE()` returns `NULL` and generates a warning:

```
mysql> SET sql_mode = '';
mysql> SELECT STR_TO_DATE('15:35:00', '%H:%i:%s');
+-----------------------------------+
| STR_TO_DATE('15:35:00', '%H:%i:%s') |
+-----------------------------------+
| 15:35:00                          |
+-----------------------------------+
mysql> SET sql_mode = 'NO_ZERO_IN_DATE';
mysql> SELECT STR_TO_DATE('15:35:00', '%h:%i:%s');
+-----------------------------------+
| STR_TO_DATE('15:35:00', '%h:%i:%s') |
+-----------------------------------+
| NULL                              |
+-----------------------------------+
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1411
Message: Incorrect datetime value: '15:35:00' for function str_to_date
```

Note

You cannot use format `"%X%V"` to convert a year-week string to a date because the combination of a year and week does not uniquely identify a year and month if the week crosses a month boundary. To convert a year-week to a date, you should also specify the weekday:

```
mysql> SELECT STR_TO_DATE('200442 Monday', '%X%V %W');
        -> '2004-10-18'
```

- `SUBDATE(`*`date`*`,INTERVAL` *`expr unit`*`)`, `SUBDATE(`*`expr`*`,`*`days`*`)`

  When invoked with the `INTERVAL` form of the second argument, `SUBDATE()` is a synonym for `DATE_SUB()`. For information on the `INTERVAL` *`unit`* argument, see the discussion for `DATE_ADD()`.

```
mysql> SELECT DATE_SUB('2008-01-02', INTERVAL 31 DAY);
        -> '2007-12-02'
mysql> SELECT SUBDATE('2008-01-02', INTERVAL 31 DAY);
        -> '2007-12-02'
```

  The second form enables the use of an integer value for *`days`*. In such cases, it is interpreted as the number of days to be subtracted from the date or datetime expression *`expr`*.

```
mysql> SELECT SUBDATE('2008-01-02 12:00:00', 31);
        -> '2007-12-02 12:00:00'
```

- `SUBTIME(`*`expr1`*`,`*`expr2`*`)`

  `SUBTIME()` returns *`expr1`* − *`expr2`* expressed as a value in the same format as *`expr1`*. *`expr1`* is a time or datetime expression, and *`expr2`* is a time expression.

```
mysql> SELECT SUBTIME('2007-12-31 23:59:59.999999','1 1:1:1.000002');
        -> '2007-12-30 22:58:58.999997'
mysql> SELECT SUBTIME('01:00:00.999999', '02:00:00.999998');
        -> '-00:59:59.999999'
```

- `SYSDATE([`*`fsp`*`])`

  Returns the current date and time as a value in `'YYYY-MM-DD HH:MM:SS'` or `YYYYMMDDHHMMSS` format, depending on whether the function is used in a string or numeric context.

If the *fsp* argument is given to specify a fractional seconds precision from 0 to 6, the return value includes a fractional seconds part of that many digits.

SYSDATE() returns the time at which it executes. This differs from the behavior for NOW(), which returns a constant time that indicates the time at which the statement began to execute. (Within a stored function or trigger, NOW() returns the time at which the function or triggering statement began to execute.)

```
mysql> SELECT NOW(), SLEEP(2), NOW();
+---------------------+----------+---------------------+
| NOW()               | SLEEP(2) | NOW()               |
+---------------------+----------+---------------------+
| 2006-04-12 13:47:36 |        0 | 2006-04-12 13:47:36 |
+---------------------+----------+---------------------+

mysql> SELECT SYSDATE(), SLEEP(2), SYSDATE();
+---------------------+----------+---------------------+
| SYSDATE()           | SLEEP(2) | SYSDATE()           |
+---------------------+----------+---------------------+
| 2006-04-12 13:47:44 |        0 | 2006-04-12 13:47:46 |
+---------------------+----------+---------------------+
```

In addition, the SET TIMESTAMP statement affects the value returned by NOW() but not by SYSDATE(). This means that timestamp settings in the binary log have no effect on invocations of SYSDATE().

Because SYSDATE() can return different values even within the same statement, and is not affected by SET TIMESTAMP, it is nondeterministic and therefore unsafe for replication if statement-based binary logging is used. If that is a problem, you can use row-based logging.

Alternatively, you can use the --sysdate-is-now option to cause SYSDATE() to be an alias for NOW(). This works if the option is used on both the master and the slave.

The nondeterministic nature of SYSDATE() also means that indexes cannot be used for evaluating expressions that refer to it.

- TIME(*expr*)

  Extracts the time part of the time or datetime expression *expr* and returns it as a string.

  This function is unsafe for statement-based replication. A warning is logged if you use this function when binlog_format is set to STATEMENT.

```
mysql> SELECT TIME('2003-12-31 01:02:03');
        -> '01:02:03'
mysql> SELECT TIME('2003-12-31 01:02:03.000123');
        -> '01:02:03.000123'
```

- TIMEDIFF(*expr1*,*expr2*)

  TIMEDIFF() returns *expr1* − *expr2* expressed as a time value. *expr1* and *expr2* are time or date-and-time expressions, but both must be of the same type.

  The result returned by TIMEDIFF() is limited to the range allowed for TIME values. Alternatively, you can use either of the functions TIMESTAMPDIFF() and UNIX_TIMESTAMP(), both of which return integers.

```
mysql> SELECT TIMEDIFF('2000:01:01 00:00:00',
    ->                 '2000:01:01 00:00:00.000001');
        -> '-00:00:00.000001'
mysql> SELECT TIMEDIFF('2008-12-31 23:59:59.000001',
    ->                 '2008-12-30 01:01:01.000002');
```

```
        -> '46:58:57.999999'
```

- `TIMESTAMP(expr)`, `TIMESTAMP(expr1,expr2)`

  With a single argument, this function returns the date or datetime expression *`expr`* as a datetime value. With two arguments, it adds the time expression *`expr2`* to the date or datetime expression *`expr1`* and returns the result as a datetime value.

```
mysql> SELECT TIMESTAMP('2003-12-31');
        -> '2003-12-31 00:00:00'
mysql> SELECT TIMESTAMP('2003-12-31 12:00:00','12:00:00');
        -> '2004-01-01 00:00:00'
```

- `TIMESTAMPADD(unit,interval,datetime_expr)`

  Adds the integer expression *`interval`* to the date or datetime expression *`datetime_expr`*. The unit for *`interval`* is given by the *`unit`* argument, which should be one of the following values: `MICROSECOND` (microseconds), `SECOND`, `MINUTE`, `HOUR`, `DAY`, `WEEK`, `MONTH`, `QUARTER`, or `YEAR`.

  The *`unit`* value may be specified using one of keywords as shown, or with a prefix of `SQL_TSI_`. For example, `DAY` and `SQL_TSI_DAY` both are legal.

```
mysql> SELECT TIMESTAMPADD(MINUTE,1,'2003-01-02');
        -> '2003-01-02 00:01:00'
mysql> SELECT TIMESTAMPADD(WEEK,1,'2003-01-02');
        -> '2003-01-09'
```

- `TIMESTAMPDIFF(unit,datetime_expr1,datetime_expr2)`

  Returns *`datetime_expr2`* − *`datetime_expr1`*, where *`datetime_expr1`* and *`datetime_expr2`* are date or datetime expressions. One expression may be a date and the other a datetime; a date value is treated as a datetime having the time part `'00:00:00'` where necessary. The unit for the result (an integer) is given by the *`unit`* argument. The legal values for *`unit`* are the same as those listed in the description of the `TIMESTAMPADD()` function.

```
mysql> SELECT TIMESTAMPDIFF(MONTH,'2003-02-01','2003-05-01');
        -> 3
mysql> SELECT TIMESTAMPDIFF(YEAR,'2002-05-01','2001-01-01');
        -> -1
mysql> SELECT TIMESTAMPDIFF(MINUTE,'2003-02-01','2003-05-01 12:05:55');
        -> 128885
```

Note

The order of the date or datetime arguments for this function is the opposite of that used with the `TIMESTAMP()` function when invoked with 2 arguments.

- `TIME_FORMAT(time,format)`

  This is used like the `DATE_FORMAT()` function, but the *`format`* string may contain format specifiers only for hours, minutes, seconds, and microseconds. Other specifiers produce a `NULL` value or `0`.

  If the *`time`* value contains an hour part that is greater than `23`, the `%H` and `%k` hour format specifiers produce a value larger than the usual range of `0..23`. The other hour format specifiers produce the hour value modulo 12.

```
mysql> SELECT TIME_FORMAT('100:00:00', '%H %k %h %I %l');
        -> '100 100 04 04 4'
```

- `TIME_TO_SEC(time)`

  Returns the *`time`* argument, converted to seconds.

```
mysql> SELECT TIME_TO_SEC('22:23:00');
        -> 80580
mysql> SELECT TIME_TO_SEC('00:39:38');
        -> 2378
```

- `TO_DAYS(date)`

  Given a date *`date`*, returns a day number (the number of days since year 0).

```
mysql> SELECT TO_DAYS(950501);
        -> 728779
mysql> SELECT TO_DAYS('2007-10-07');
        -> 733321
```

  `TO_DAYS()` is not intended for use with values that precede the advent of the Gregorian calendar (1582), because it does not take into account the days that were lost when the calendar was changed. For dates before 1582 (and possibly a later year in other locales), results from this function are not reliable. See Section 12.8, “What Calendar Is Used By MySQL?”, for details.

  Remember that MySQL converts two-digit year values in dates to four-digit form using the rules in Section 11.3, “Date and Time Types”. For example, `'2008-10-07'` and `'08-10-07'` are seen as identical dates:

```
mysql> SELECT TO_DAYS('2008-10-07'), TO_DAYS('08-10-07');
        -> 733687, 733687
```

  In MySQL, the zero date is defined as `'0000-00-00'`, even though this date is itself considered invalid. This means that, for `'0000-00-00'` and `'0000-01-01'`, `TO_DAYS()` returns the values shown here:

```
mysql> SELECT TO_DAYS('0000-00-00');
+-----------------------+
| to_days('0000-00-00') |
+-----------------------+
|                  NULL |
+-----------------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS;
+---------+------+----------------------------------------+
| Level   | Code | Message                                |
+---------+------+----------------------------------------+
| Warning | 1292 | Incorrect datetime value: '0000-00-00' |
+---------+------+----------------------------------------+
1 row in set (0.00 sec)


mysql> SELECT TO_DAYS('0000-01-01');
+-----------------------+
| to_days('0000-01-01') |
+-----------------------+
|                     1 |
+-----------------------+
1 row in set (0.00 sec)
```

  This is true whether or not the `ALLOW_INVALID_DATES` SQL server mode is enabled.

- `TO_SECONDS(expr)`

  Given a date or datetime *`expr`*, returns the number of seconds since the year 0. If *`expr`* is not a valid date or datetime value, returns `NULL`.

```
```

```
mysql> SELECT TO_SECONDS(950501);
        -> 62966505600
mysql> SELECT TO_SECONDS('2009-11-29');
        -> 63426672000
mysql> SELECT TO_SECONDS('2009-11-29 13:43:32');
        -> 63426721412
mysql> SELECT TO_SECONDS( NOW() );
        -> 63426721458
```

Like `TO_DAYS()`, `TO_SECONDS()` is not intended for use with values that precede the advent of the Gregorian calendar (1582), because it does not take into account the days that were lost when the calendar was changed. For dates before 1582 (and possibly a later year in other locales), results from this function are not reliable. See Section 12.8, “What Calendar Is Used By MySQL?”, for details.

Like `TO_DAYS()`, `TO_SECONDS()`, converts two-digit year values in dates to four-digit form using the rules in Section 11.3, “Date and Time Types”.

In MySQL, the zero date is defined as `'0000-00-00'`, even though this date is itself considered invalid. This means that, for `'0000-00-00'` and `'0000-01-01'`, `TO_SECONDS()` returns the values shown here:

```
mysql> SELECT TO_SECONDS('0000-00-00');
+--------------------------+
| TO_SECONDS('0000-00-00') |
+--------------------------+
|                     NULL |
+--------------------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS;
+---------+------+---------------------------------------+
| Level   | Code | Message                               |
+---------+------+---------------------------------------+
| Warning | 1292 | Incorrect datetime value: '0000-00-00' |
+---------+------+---------------------------------------+
1 row in set (0.00 sec)


mysql> SELECT TO_SECONDS('0000-01-01');
+--------------------------+
| TO_SECONDS('0000-01-01') |
+--------------------------+
|                    86400 |
+--------------------------+
1 row in set (0.00 sec)
```

This is true whether or not the `ALLOW_INVALID_DATES` SQL server mode is enabled.

- `UNIX_TIMESTAMP()`, `UNIX_TIMESTAMP(`*`date`*`)`

  If called with no argument, returns a Unix timestamp (seconds since `'1970-01-01 00:00:00'` UTC). The return value is an integer if no argument is given or the argument does not include a fractional seconds part, or `DECIMAL` if an argument is given that includes a fractional seconds part.

  If `UNIX_TIMESTAMP()` is called with a *`date`* argument, it returns the value of the argument as seconds since `'1970-01-01 00:00:00'` UTC. The *`date`* argument may be a `DATE`, `DATETIME`, or `TIMESTAMP` string, or a number in `YYMMDD`, `YYMMDDHHMMSS`, `YYYYMMDD`, or `YYYYMMDDHHMMSS` format. If the argument includes a time part, it may optionally include a fractional seconds part. The server interprets *`date`* as a value in the current time zone and converts it to an internal value in UTC. Clients can set their time zone as described in Section 10.6, “MySQL Server Time Zone Support”.

  ```
  mysql> SELECT UNIX_TIMESTAMP();
          -> 1447431666
  mysql> SELECT UNIX_TIMESTAMP('2015-11-13 10:20:19');
  ```

```
        -> 1447431619
mysql> SELECT UNIX_TIMESTAMP('2015-11-13 10:20:19.012');
        -> 1447431619.012
```

When `UNIX_TIMESTAMP()` is used on a `TIMESTAMP` column, the function returns the internal timestamp value directly, with no implicit “string-to-Unix-timestamp” conversion. If you pass an out-of-range date to `UNIX_TIMESTAMP()`, it returns `0`. The valid range of values is the same as for the `TIMESTAMP` data type: `'1970-01-01 00:00:01.000000'` UTC to `'2038-01-19 03:14:07.999999'` UTC.

Note: If you use `UNIX_TIMESTAMP()` and `FROM_UNIXTIME()` to convert between `TIMESTAMP` values and Unix timestamp values, the conversion is lossy because the mapping is not one-to-one in both directions. For example, due to conventions for local time zone changes, it is possible for two `UNIX_TIMESTAMP()` to map two `TIMESTAMP` values to the same Unix timestamp value. `FROM_UNIXTIME()` will map that value back to only one of the original `TIMESTAMP` values. Here is an example, using `TIMESTAMP` values in the `CET` time zone:

```
mysql> SELECT UNIX_TIMESTAMP('2005-03-27 03:00:00');
+---------------------------------------+
| UNIX_TIMESTAMP('2005-03-27 03:00:00') |
+---------------------------------------+
|                            1111885200 |
+---------------------------------------+
mysql> SELECT UNIX_TIMESTAMP('2005-03-27 02:00:00');
+---------------------------------------+
| UNIX_TIMESTAMP('2005-03-27 02:00:00') |
+---------------------------------------+
|                            1111885200 |
+---------------------------------------+
mysql> SELECT FROM_UNIXTIME(1111885200);
+---------------------------+
| FROM_UNIXTIME(1111885200) |
+---------------------------+
| 2005-03-27 03:00:00       |
+---------------------------+
```

If you want to subtract `UNIX_TIMESTAMP()` columns, you might want to cast the result to signed integers. See Section 12.10, “Cast Functions and Operators”.

- `UTC_DATE`, `UTC_DATE()`

  Returns the current UTC date as a value in `'YYYY-MM-DD'` or `YYYYMMDD` format, depending on whether the function is used in a string or numeric context.

```
mysql> SELECT UTC_DATE(), UTC_DATE() + 0;
        -> '2003-08-14', 20030814
```

- `UTC_TIME`, `UTC_TIME([fsp])`

  Returns the current UTC time as a value in `'HH:MM:SS'` or `HHMMSS` format, depending on whether the function is used in a string or numeric context.

  If the *`fsp`* argument is given to specify a fractional seconds precision from 0 to 6, the return value includes a fractional seconds part of that many digits.

```
mysql> SELECT UTC_TIME(), UTC_TIME() + 0;
        -> '18:07:53', 180753.000000
```

- `UTC_TIMESTAMP`, `UTC_TIMESTAMP([fsp])`

  Returns the current UTC date and time as a value in `'YYYY-MM-DD HH:MM:SS'` or `YYYYMMDDHHMMSS` format, depending on whether the function is used in a string or numeric context.

If the *fsp* argument is given to specify a fractional seconds precision from 0 to 6, the return value includes a fractional seconds part of that many digits.

```
mysql> SELECT UTC_TIMESTAMP(), UTC_TIMESTAMP() + 0;
        -> '2003-08-14 18:08:04', 20030814180804.000000
```

- `WEEK(date[,mode])`

  This function returns the week number for *date*. The two-argument form of `WEEK()` enables you to specify whether the week starts on Sunday or Monday and whether the return value should be in the range from `0` to `53` or from `1` to `53`. If the *mode* argument is omitted, the value of the `default_week_format` system variable is used. See Section 5.1.5, “Server System Variables”.

  The following table describes how the *mode* argument works.

Mode	First day of week	Range	Week 1 is the first week …
0	Sunday	0-53	with a Sunday in this year
1	Monday	0-53	with 4 or more days this year
2	Sunday	1-53	with a Sunday in this year
3	Monday	1-53	with 4 or more days this year
4	Sunday	0-53	with 4 or more days this year
5	Monday	0-53	with a Monday in this year
6	Sunday	1-53	with 4 or more days this year
7	Monday	1-53	with a Monday in this year

  For *mode* values with a meaning of “with 4 or more days this year,” weeks are numbered according to ISO 8601:1988:

  - If the week containing January 1 has 4 or more days in the new year, it is week 1.
  - Otherwise, it is the last week of the previous year, and the next week is week 1.

```
mysql> SELECT WEEK('2008-02-20');
        -> 7
mysql> SELECT WEEK('2008-02-20',0);
        -> 7
mysql> SELECT WEEK('2008-02-20',1);
        -> 8
mysql> SELECT WEEK('2008-12-31',1);
        -> 53
```

  Note that if a date falls in the last week of the previous year, MySQL returns `0` if you do not use `2`, `3`, `6`, or `7` as the optional *mode* argument:

```
mysql> SELECT YEAR('2000-01-01'), WEEK('2000-01-01',0);
        -> 2000, 0
```

  One might argue that `WEEK()` should return `52` because the given date actually occurs in the 52nd week of 1999. `WEEK()` returns `0` instead so that the return value is “the week number in the given year.” This makes use of the `WEEK()` function reliable when combined with other functions that extract a date part from a date.

  If you prefer a result evaluated with respect to the year that contains the first day of the week for the given date, use `0`, `2`, `5`, or `7` as the optional *mode* argument.

```
mysql> SELECT WEEK('2000-01-01',2);
        -> 52
```

Alternatively, use the `YEARWEEK()` function:

```
mysql> SELECT YEARWEEK('2000-01-01');
        -> 199952
mysql> SELECT MID(YEARWEEK('2000-01-01'),5,2);
        -> '52'
```

- `WEEKDAY(date)`

  Returns the weekday index for *`date`* (`0` = Monday, `1` = Tuesday, … `6` = Sunday).

```
mysql> SELECT WEEKDAY('2008-02-03 22:23:00');
        -> 6
mysql> SELECT WEEKDAY('2007-11-06');
        -> 1
```

- `WEEKOFYEAR(date)`

  Returns the calendar week of the date as a number in the range from `1` to `53`. `WEEKOFYEAR()` is a compatibility function that is equivalent to `WEEK(date,3)`.

```
mysql> SELECT WEEKOFYEAR('2008-02-20');
        -> 8
```

- `YEAR(date)`

  Returns the year for *`date`*, in the range `1000` to `9999`, or `0` for the “zero” date.

```
mysql> SELECT YEAR('1987-01-01');
        -> 1987
```

- `YEARWEEK(date)`, `YEARWEEK(date,mode)`

  Returns year and week for a date. The year in the result may be different from the year in the date argument for the first and the last week of the year.

  The *`mode`* argument works exactly like the *`mode`* argument to `WEEK()`. For the single-argument syntax, a *`mode`* value of 0 is used. Unlike `WEEK()`, the value of `default_week_format` does not influence `YEARWEEK()`.

```
mysql> SELECT YEARWEEK('1987-01-01');
        -> 198652
```

  Note that the week number is different from what the `WEEK()` function would return (`0`) for optional arguments `0` or `1`, as `WEEK()` then returns the week in the context of the given year.

12.8 What Calendar Is Used By MySQL?

MySQL uses what is known as a *proleptic Gregorian calendar*.

Every country that has switched from the Julian to the Gregorian calendar has had to discard at least ten days during the switch. To see how this works, consider the month of October 1582, when the first Julian-to-Gregorian switch occurred.

Monday	Tuesday	Wednesday	Thursday	Friday	Saturday	Sunday
1	2	3	4	15	16	17
18	19	20	21	22	23	24

Monday	Tuesday	Wednesday	Thursday	Friday	Saturday	Sunday
25	26	27	28	29	30	31

There are no dates between October 4 and October 15. This discontinuity is called the *cutover*. Any dates before the cutover are Julian, and any dates following the cutover are Gregorian. Dates during a cutover are nonexistent.

A calendar applied to dates when it was not actually in use is called *proleptic*. Thus, if we assume there was never a cutover and Gregorian rules always rule, we have a proleptic Gregorian calendar. This is what is used by MySQL, as is required by standard SQL. For this reason, dates prior to the cutover stored as MySQL `DATE` or `DATETIME` values must be adjusted to compensate for the difference. It is important to realize that the cutover did not occur at the same time in all countries, and that the later it happened, the more days were lost. For example, in Great Britain, it took place in 1752, when Wednesday September 2 was followed by Thursday September 14. Russia remained on the Julian calendar until 1918, losing 13 days in the process, and what is popularly referred to as its “October Revolution” occurred in November according to the Gregorian calendar.

12.9 Full-Text Search Functions

```
MATCH (col1,col2,...) AGAINST (expr [search_modifier])
```

```
search_modifier:
  {
       IN NATURAL LANGUAGE MODE
     | IN NATURAL LANGUAGE MODE WITH QUERY EXPANSION
     | IN BOOLEAN MODE
     | WITH QUERY EXPANSION
  }
```

MySQL has support for full-text indexing and searching:

- A full-text index in MySQL is an index of type `FULLTEXT`.
- Full-text indexes can be used only with `InnoDB` or `MyISAM` tables, and can be created only for `CHAR`, `VARCHAR`, or `TEXT` columns.
- `MySQL` provides a built-in full-text ngram parser that supports Chinese, Japanese, and Korean (CJK), and an installable MeCab full-text parser plugin for Japanese. Parsing differences are outlined in Section 12.9.8, “ngram Full-Text Parser”, and Section 12.9.9, “MeCab Full-Text Parser Plugin”.
- A `FULLTEXT` index definition can be given in the `CREATE TABLE` statement when a table is created, or added later using `ALTER TABLE` or `CREATE INDEX`.
- For large data sets, it is much faster to load your data into a table that has no `FULLTEXT` index and then create the index after that, than to load data into a table that has an existing `FULLTEXT` index.

Full-text searching is performed using `MATCH() ... AGAINST` syntax. `MATCH()` takes a comma-separated list that names the columns to be searched. `AGAINST` takes a string to search for, and an optional modifier that indicates what type of search to perform. The search string must be a string value that is constant during query evaluation. This rules out, for example, a table column because that can differ for each row.

There are three types of full-text searches:

- A natural language search interprets the search string as a phrase in natural human language (a phrase in free text). There are no special operators, with the exception of double quote (") characters. The stopword list applies. For more information about stopword lists, see Section 12.9.4, “Full-Text Stopwords”.

  Full-text searches are natural language searches if the `IN NATURAL LANGUAGE MODE` modifier is given or if no modifier is given. For more information, see Section 12.9.1, “Natural Language Full-Text Searches”.

- A boolean search interprets the search string using the rules of a special query language. The string contains the words to search for. It can also contain operators that specify requirements such that a word must be present or absent in matching rows, or that it should be weighted higher or lower than usual. Certain common words (stopwords) are omitted from the search index and do not match if present in the search string. The `IN BOOLEAN MODE` modifier specifies a boolean search. For more information, see Section 12.9.2, “Boolean Full-Text Searches”.

- A query expansion search is a modification of a natural language search. The search string is used to perform a natural language search. Then words from the most relevant rows returned by the search are added to the search string and the search is done again. The query returns the rows from the second search. The `IN NATURAL LANGUAGE MODE WITH QUERY EXPANSION` or `WITH QUERY EXPANSION` modifier specifies a query expansion search. For more information, see Section 12.9.3, “Full-Text Searches with Query Expansion”.

For information about `FULLTEXT` query performance, see Section 8.3.4, “Column Indexes”.

For more information about `InnoDB FULLTEXT` indexes, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”.

Constraints on full-text searching are listed in Section 12.9.5, “Full-Text Restrictions”.

The `myisam_ftdump` utility dumps the contents of a `MyISAM` full-text index. This may be helpful for debugging full-text queries. See Section 4.6.3, “`myisam_ftdump` — Display Full-Text Index information”.

12.9.1 Natural Language Full-Text Searches

By default or with the `IN NATURAL LANGUAGE MODE` modifier, the `MATCH()` function performs a natural language search for a string against a *text collection*. A collection is a set of one or more columns included in a `FULLTEXT` index. The search string is given as the argument to `AGAINST()`. For each row in the table, `MATCH()` returns a relevance value; that is, a similarity measure between the search string and the text in that row in the columns named in the `MATCH()` list.

```
mysql> CREATE TABLE articles (
          id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
          title VARCHAR(200),
          body TEXT,
          FULLTEXT (title,body)
        ) ENGINE=InnoDB;
Query OK, 0 rows affected (0.08 sec)

mysql> INSERT INTO articles (title,body) VALUES
        ('MySQL Tutorial','DBMS stands for DataBase ...'),
        ('How To Use MySQL Well','After you went through a ...'),
        ('Optimizing MySQL','In this tutorial we will show ...'),
        ('1001 MySQL Tricks','1. Never run mysqld as root. 2. ...'),
        ('MySQL vs. YourSQL','In the following database comparison ...'),
        ('MySQL Security','When configured properly, MySQL ...');
Query OK, 6 rows affected (0.01 sec)
Records: 6  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM articles
        WHERE MATCH (title,body)
        AGAINST ('database' IN NATURAL LANGUAGE MODE);
+----+-------------------+------------------------------------------+
| id | title             | body                                     |
+----+-------------------+------------------------------------------+
|  1 | MySQL Tutorial    | DBMS stands for DataBase ...             |
|  5 | MySQL vs. YourSQL | In the following database comparison ... |
+----+-------------------+------------------------------------------+
2 rows in set (0.00 sec)
```

By default, the search is performed in case-insensitive fashion. To perform a case-sensitive full-text search, use a case-sensitive or binary collation for the indexed columns. For example, a column

that uses the utf8mb4 character set of can be assigned a collation of utf8mb4_0900_as_cs or utf8mb4_bin to make it case sensitive for full-text searches.

When MATCH() is used in a WHERE clause, as in the example shown earlier, the rows returned are automatically sorted with the highest relevance first. Relevance values are nonnegative floating-point numbers. Zero relevance means no similarity. Relevance is computed based on the number of words in the row (document), the number of unique words in the row, the total number of words in the collection, and the number of rows that contain a particular word.

Note

The term “document” may be used interchangeably with the term “row”, and both terms refer to the indexed part of the row. The term “collection” refers to the indexed columns and encompasses all rows.

To simply count matches, you could use a query like this:

```
mysql> SELECT COUNT(*) FROM articles
    WHERE MATCH (title,body)
    AGAINST ('database' IN NATURAL LANGUAGE MODE);
+----------+
| COUNT(*) |
+----------+
|        2 |
+----------+
1 row in set (0.00 sec)
```

You might find it quicker to rewrite the query as follows:

```
mysql> SELECT
    COUNT(IF(MATCH (title,body) AGAINST ('database' IN NATURAL LANGUAGE MODE), 1, NULL))
    AS count
    FROM articles;
+-------+
| count |
+-------+
|     2 |
+-------+
1 row in set (0.03 sec)
```

The first query does some extra work (sorting the results by relevance) but also can use an index lookup based on the WHERE clause. The index lookup might make the first query faster if the search matches few rows. The second query performs a full table scan, which might be faster than the index lookup if the search term was present in most rows.

For natural-language full-text searches, the columns named in the MATCH() function must be the same columns included in some FULLTEXT index in your table. For the preceding query, note that the columns named in the MATCH() function (title and body) are the same as those named in the definition of the article table's FULLTEXT index. To search the title or body separately, you would create separate FULLTEXT indexes for each column.

You can also perform a boolean search or a search with query expansion. These search types are described in Section 12.9.2, “Boolean Full-Text Searches”, and Section 12.9.3, “Full-Text Searches with Query Expansion”.

A full-text search that uses an index can name columns only from a single table in the MATCH() clause because an index cannot span multiple tables. For MyISAM tables, a boolean search can be done in the absence of an index (albeit more slowly), in which case it is possible to name columns from multiple tables.

The preceding example is a basic illustration that shows how to use the MATCH() function where rows are returned in order of decreasing relevance. The next example shows how to retrieve the relevance values explicitly. Returned rows are not ordered because the SELECT statement includes neither WHERE nor ORDER BY clauses:

```
mysql> SELECT id, MATCH (title,body)
    AGAINST ('Tutorial' IN NATURAL LANGUAGE MODE) AS score
    FROM articles;
+----+---------------------+
| id | score               |
+----+---------------------+
|  1 | 0.22764469683170319 |
|  2 |                   0 |
|  3 | 0.22764469683170319 |
|  4 |                   0 |
|  5 |                   0 |
|  6 |                   0 |
+----+---------------------+
6 rows in set (0.00 sec)
```

The following example is more complex. The query returns the relevance values and it also sorts the rows in order of decreasing relevance. To achieve this result, specify `MATCH()` twice: once in the `SELECT` list and once in the `WHERE` clause. This causes no additional overhead, because the MySQL optimizer notices that the two `MATCH()` calls are identical and invokes the full-text search code only once.

```
mysql> SELECT id, body, MATCH (title,body) AGAINST
    ('Security implications of running MySQL as root'
    IN NATURAL LANGUAGE MODE) AS score
    FROM articles WHERE MATCH (title,body) AGAINST
    ('Security implications of running MySQL as root'
    IN NATURAL LANGUAGE MODE);
+----+-------------------------------------+-----------------+
| id | body                                | score           |
+----+-------------------------------------+-----------------+
|  4 | 1. Never run mysqld as root. 2. ... | 1.5219271183014 |
|  6 | When configured properly, MySQL ... | 1.3114095926285 |
+----+-------------------------------------+-----------------+
2 rows in set (0.00 sec)
```

A phrase that is enclosed within double quote (`"`) characters matches only rows that contain the phrase *literally, as it was typed*. The full-text engine splits the phrase into words and performs a search in the `FULLTEXT` index for the words. Nonword characters need not be matched exactly: Phrase searching requires only that matches contain exactly the same words as the phrase and in the same order. For example, `"test phrase"` matches `"test, phrase"`. If the phrase contains no words that are in the index, the result is empty. For example, if all words are either stopwords or shorter than the minimum length of indexed words, the result is empty.

The MySQL `FULLTEXT` implementation regards any sequence of true word characters (letters, digits, and underscores) as a word. That sequence may also contain apostrophes (`'`), but not more than one in a row. This means that `aaa'bbb` is regarded as one word, but `aaa''bbb` is regarded as two words. Apostrophes at the beginning or the end of a word are stripped by the `FULLTEXT` parser; `'aaa'bbb'` would be parsed as `aaa'bbb`.

The built-in `FULLTEXT` parser determines where words start and end by looking for certain delimiter characters; for example, (space), `,` (comma), and `.` (period). If words are not separated by delimiters (as in, for example, Chinese), the built-in `FULLTEXT` parser cannot determine where a word begins or ends. To be able to add words or other indexed terms in such languages to a `FULLTEXT` index that uses the built-in `FULLTEXT` parser, you must preprocess them so that they are separated by some arbitrary delimiter. Alternatively, you can create `FULLTEXT` indexes using the ngram parser plugin (for Chinese, Japanese, or Korean) or the MeCab parser plugin (for Japanese).

It is possible to write a plugin that replaces the built-in full-text parser. For details, see Section 28.2, “The MySQL Plugin API”. For example parser plugin source code, see the `plugin/fulltext` directory of a MySQL source distribution.

Some words are ignored in full-text searches:

- Any word that is too short is ignored. The default minimum length of words that are found by full-text searches is three characters for `InnoDB` search indexes, or four characters for `MyISAM`. You can control the cutoff by setting a configuration option before creating the index: `innodb_ft_min_token_size` configuration option for `InnoDB` search indexes, or `ft_min_word_len` for `MyISAM`.

  Note

  This behavior does not apply to `FULLTEXT` indexes that use the ngram parser. For the ngram parser, token length is defined by the `ngram_token_size` option.

- Words in the stopword list are ignored. A stopword is a word such as “the” or “some” that is so common that it is considered to have zero semantic value. There is a built-in stopword list, but it can be overridden by a user-defined list. The stopword lists and related configuration options are different for `InnoDB` search indexes and `MyISAM` ones. Stopword processing is controlled by the configuration options `innodb_ft_enable_stopword`, `innodb_ft_server_stopword_table`, and `innodb_ft_user_stopword_table` for `InnoDB` search indexes, and `ft_stopword_file` for `MyISAM` ones.

See Section 12.9.4, “Full-Text Stopwords” to view default stopword lists and how to change them. The default minimum word length can be changed as described in Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.

Every correct word in the collection and in the query is weighted according to its significance in the collection or query. Thus, a word that is present in many documents has a lower weight, because it has lower semantic value in this particular collection. Conversely, if the word is rare, it receives a higher weight. The weights of the words are combined to compute the relevance of the row. This technique works best with large collections.

MyISAM Limitation

For very small tables, word distribution does not adequately reflect their semantic value, and this model may sometimes produce bizarre results for search indexes on `MyISAM` tables. For example, although the word “MySQL” is present in every row of the `articles` table shown earlier, a search for the word in a `MyISAM` search index produces no results:

```
mysql> SELECT * FROM articles
    WHERE MATCH (title,body)
    AGAINST ('MySQL' IN NATURAL LANGUAGE MODE);
Empty set (0.00 sec)
```

The search result is empty because the word “MySQL” is present in at least 50% of the rows, and so is effectively treated as a stopword. This filtering technique is more suitable for large data sets, where you might not want the result set to return every second row from a 1GB table, than for small data sets where it might cause poor results for popular terms.

The 50% threshold can surprise you when you first try full-text searching to see how it works, and makes `InnoDB` tables more suited to experimentation with full-text searches. If you create a `MyISAM` table and insert only one or two rows of text into it, every word in the text occurs in at least 50% of the rows. As a result, no search returns any results until the table contains more rows. Users who need to bypass the 50% limitation can build search indexes on `InnoDB` tables, or use the boolean search mode explained in Section 12.9.2, “Boolean Full-Text Searches”.

12.9.2 Boolean Full-Text Searches

MySQL can perform boolean full-text searches using the `IN BOOLEAN MODE` modifier. With this modifier, certain characters have special meaning at the beginning or end of words in the search string. In the following query, the `+` and `-` operators indicate that a word must be present or absent, respectively, for a match to occur. Thus, the query retrieves all the rows that contain the word “MySQL” but that do *not* contain the word “YourSQL”:

```
mysql> SELECT * FROM articles WHERE MATCH (title,body)
    AGAINST ('+MySQL -YourSQL' IN BOOLEAN MODE);
+----+-----------------------+-------------------------------------+
| id | title                 | body                                |
+----+-----------------------+-------------------------------------+
|  1 | MySQL Tutorial        | DBMS stands for DataBase ...        |
|  2 | How To Use MySQL Well | After you went through a ...        |
|  3 | Optimizing MySQL      | In this tutorial we will show ...   |
|  4 | 1001 MySQL Tricks     | 1. Never run mysqld as root. 2. ... |
|  6 | MySQL Security        | When configured properly, MySQL ... |
+----+-----------------------+-------------------------------------+
```

Note

In implementing this feature, MySQL uses what is sometimes referred to as *implied Boolean logic*, in which

- `+` stands for `AND`
- `-` stands for `NOT`
- [*no operator*] implies `OR`

Boolean full-text searches have these characteristics:

- They do not automatically sort rows in order of decreasing relevance.
- `InnoDB` tables require a `FULLTEXT` index on all columns of the `MATCH()` expression to perform boolean queries. Boolean queries against a `MyISAM` search index can work even without a `FULLTEXT` index, although a search executed in this fashion would be quite slow.
- The minimum and maximum word length full-text parameters apply to `FULLTEXT` indexes created using the built-in `FULLTEXT` parser and MeCab parser plugin. `innodb_ft_min_token_size` and `innodb_ft_max_token_size` are used for `InnoDB` search indexes. `ft_min_word_len` and `ft_max_word_len` are used for `MyISAM` search indexes.

  Minimum and maximum word length full-text parameters do not apply to `FULLTEXT` indexes created using the ngram parser. ngram token size is defined by the `ngram_token_size` option.
- The stopword list applies, controlled by `innodb_ft_enable_stopword`, `innodb_ft_server_stopword_table`, and `innodb_ft_user_stopword_table` for `InnoDB` search indexes, and `ft_stopword_file` for `MyISAM` ones.
- `InnoDB` full-text search does not support the use of multiple operators on a single search word, as in this example: `'++apple'`. Use of multiple operators on a single search word returns a syntax error to standard out. MyISAM full-text search will successfully process the same search ignoring all operators except for the operator immediately adjacent to the search word.
- `InnoDB` full-text search only supports leading plus or minus signs. For example, `InnoDB` supports `'+apple'` but does not support `'apple+'`. Specifying a trailing plus or minus sign causes `InnoDB` to report a syntax error.
- `InnoDB` full-text search does not support the use of a leading plus sign with wildcard (`'+*'`), a plus and minus sign combination (`'+-'`), or leading a plus and minus sign combination (`'+-apple'`). These invalid queries return a syntax error.
- `InnoDB` full-text search does not support the use of the `@` symbol in boolean full-text searches. The `@` symbol is reserved for use by the `@distance` proximity search operator.

- They do not use the 50% threshold that applies to `MyISAM` search indexes.

The boolean full-text search capability supports the following operators:

- `+`

  A leading or trailing plus sign indicates that this word *must* be present in each row that is returned. `InnoDB` only supports leading plus signs.

- `-`

  A leading or trailing minus sign indicates that this word must *not* be present in any of the rows that are returned. `InnoDB` only supports leading minus signs.

  Note: The `-` operator acts only to exclude rows that are otherwise matched by other search terms. Thus, a boolean-mode search that contains only terms preceded by `-` returns an empty result. It does not return “all rows except those containing any of the excluded terms.”

- (no operator)

  By default (when neither `+` nor `-` is specified), the word is optional, but the rows that contain it are rated higher. This mimics the behavior of `MATCH() ... AGAINST()` without the `IN BOOLEAN MODE` modifier.

- `@distance`

  This operator works on `InnoDB` tables only. It tests whether two or more words all start within a specified distance from each other, measured in words. Specify the search words within a double-quoted string immediately before the `@distance` operator, for example, `MATCH(col1) AGAINST('"word1 word2 word3" @8' IN BOOLEAN MODE)`

- `> <`

  These two operators are used to change a word's contribution to the relevance value that is assigned to a row. The `>` operator increases the contribution and the `<` operator decreases it. See the example following this list.

- `( )`

  Parentheses group words into subexpressions. Parenthesized groups can be nested.

- `~`

  A leading tilde acts as a negation operator, causing the word's contribution to the row's relevance to be negative. This is useful for marking “noise” words. A row containing such a word is rated lower than others, but is not excluded altogether, as it would be with the `-` operator.

- `*`

  The asterisk serves as the truncation (or wildcard) operator. Unlike the other operators, it is *appended* to the word to be affected. Words match if they begin with the word preceding the `*` operator.

  If a word is specified with the truncation operator, it is not stripped from a boolean query, even if it is too short or a stopword. Whether a word is too short is determined from the `innodb_ft_min_token_size` setting for `InnoDB` tables, or `ft_min_word_len` for `MyISAM` tables. These options are not applicable to `FULLTEXT` indexes that use the ngram parser.

  The wildcarded word is considered as a prefix that must be present at the start of one or more words. If the minimum word length is 4, a search for `'+word +the*'` could return fewer rows than a search for `'+word +the'`, because the second query ignores the too-short search term `the`.

- `"`

A phrase that is enclosed within double quote (") characters matches only rows that contain the phrase *literally, as it was typed*. The full-text engine splits the phrase into words and performs a search in the `FULLTEXT` index for the words. Nonword characters need not be matched exactly: Phrase searching requires only that matches contain exactly the same words as the phrase and in the same order. For example, `"test phrase"` matches `"test, phrase"`.

If the phrase contains no words that are in the index, the result is empty. The words might not be in the index because of a combination of factors: if they do not exist in the text, are stopwords, or are shorter than the minimum length of indexed words.

The following examples demonstrate some search strings that use boolean full-text operators:

- `'apple banana'`

  Find rows that contain at least one of the two words.

- `'+apple +juice'`

  Find rows that contain both words.

- `'+apple macintosh'`

  Find rows that contain the word “apple”, but rank rows higher if they also contain “macintosh”.

- `'+apple -macintosh'`

  Find rows that contain the word “apple” but not “macintosh”.

- `'+apple ~macintosh'`

  Find rows that contain the word “apple”, but if the row also contains the word “macintosh”, rate it lower than if row does not. This is “softer” than a search for `'+apple -macintosh'`, for which the presence of “macintosh” causes the row not to be returned at all.

- `'+apple +(>turnover <strudel)'`

  Find rows that contain the words “apple” and “turnover”, or “apple” and “strudel” (in any order), but rank “apple turnover” higher than “apple strudel”.

- `'apple*'`

  Find rows that contain words such as “apple”, “apples”, “applesauce”, or “applet”.

- `'"some words"'`

  Find rows that contain the exact phrase “some words” (for example, rows that contain “some words of wisdom” but not “some noise words”). Note that the `"` characters that enclose the phrase are operator characters that delimit the phrase. They are not the quotation marks that enclose the search string itself.

Relevancy Rankings for InnoDB Boolean Mode Search

`InnoDB` full-text search is modeled on the Sphinx full-text search engine, and the algorithms used are based on BM25 and TF-IDF ranking algorithms. For these reasons, relevancy rankings for `InnoDB` boolean full-text search may differ from `MyISAM` relevancy rankings.

`InnoDB` uses a variation of the “term frequency-inverse document frequency” (`TF-IDF`) weighting system to rank a document's relevance for a given full-text search query. The `TF-IDF` weighting is based on how frequently a word appears in a document, offset by how frequently the word appears in all documents in the collection. In other words, the more frequently a word appears in a document, and the less frequently the word appears in the document collection, the higher the document is ranked.

How Relevancy Ranking is Calculated

The term frequency (`TF`) value is the number of times that a word appears in a document. The inverse document frequency (`IDF`) value of a word is calculated using the following formula, where `total_records` is the number of records in the collection, and `matching_records` is the number of records that the search term appears in.

```
${IDF} = log10( ${total_records} / ${matching_records} )
```

When a document contains a word multiple times, the IDF value is multiplied by the TF value:

```
${TF} * ${IDF}
```

Using the `TF` and `IDF` values, the relevancy ranking for a document is calculated using this formula:

```
${rank} = ${TF} * ${IDF} * ${IDF}
```

The formula is demonstrated in the following examples.

Relevancy Ranking for a Single Word Search

This example demonstrates the relevancy ranking calculation for a single-word search.

```
mysql> CREATE TABLE articles (
id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
title VARCHAR(200),
body TEXT,
FULLTEXT (title,body)
) ENGINE=InnoDB;
Query OK, 0 rows affected (1.04 sec)

mysql> INSERT INTO articles (title,body) VALUES
('MySQL Tutorial','This database tutorial ...'),
("How To Use MySQL",'After you went through a ...'),
('Optimizing Your Database','In this database tutorial ...'),
('MySQL vs. YourSQL','When comparing databases ...'),
('MySQL Security','When configured properly, MySQL ...'),
('Database, Database, Database','database database database'),
('1001 MySQL Tricks','1. Never run mysqld as root. 2. ...'),
('MySQL Full-Text Indexes', 'MySQL fulltext indexes use a ..');
Query OK, 8 rows affected (0.06 sec)
Records: 8  Duplicates: 0  Warnings: 0

mysql> SELECT id, title, body, MATCH (title,body)  AGAINST ('database' IN BOOLEAN MODE)
AS score FROM articles ORDER BY score DESC;
+----+-----------------------------+-------------------------------------+----------------------+
| id | title                       | body                                | score                |
+----+-----------------------------+-------------------------------------+----------------------+
|  6 | Database, Database, Database | database database database          |  1.0886961221694946 |
|  3 | Optimizing Your Database     | In this database tutorial ...       | 0.36289870738983154 |
|  1 | MySQL Tutorial               | This database tutorial ...          | 0.18144935369491577 |
|  2 | How To Use MySQL             | After you went through a ...        |                   0 |
|  4 | MySQL vs. YourSQL            | When comparing databases ...        |                   0 |
|  5 | MySQL Security               | When configured properly, MySQL ... |                   0 |
|  7 | 1001 MySQL Tricks            | 1. Never run mysqld as root. 2. ... |                   0 |
|  8 | MySQL Full-Text Indexes      | MySQL fulltext indexes use a ..     |                   0 |
+----+-----------------------------+-------------------------------------+----------------------+
8 rows in set (0.00 sec)
```

There are 8 records in total, with 3 that match the “database” search term. The first record (`id 6`) contains the search term 6 times and has a relevancy ranking of `1.0886961221694946`. This ranking value is calculated using a `TF` value of 6 (the “database” search term appears 6 times in record `id 6`) and an `IDF` value of 0.42596873216370745, which is calculated as follows (where 8 is the total number of records and 3 is the number of records that the search term appears in):

```
${IDF} = log10( 8 / 3 ) = 0.42596873216370745
```

The `TF` and `IDF` values are then entered into the ranking formula:

```
${rank} = ${TF} * ${IDF} * ${IDF}
```

Performing the calculation in the MySQL command-line client returns a ranking value of 1.088696164686938.

```
mysql> SELECT 6*log10(8/3)*log10(8/3);
+-------------------------+
| 6*log10(8/3)*log10(8/3) |
+-------------------------+
|       1.088696164686938 |
+-------------------------+
1 row in set (0.00 sec)
```

Note

You may notice a slight difference in the ranking values returned by the `SELECT ... MATCH ... AGAINST` statement and the MySQL command-line client (`1.0886961221694946` versus `1.088696164686938`). The difference is due to how the casts between integers and floats/doubles are performed internally by `InnoDB` (along with related precision and rounding decisions), and how they are performed elsewhere, such as in the MySQL command-line client or other types of calculators.

Relevancy Ranking for a Multiple Word Search

This example demonstrates the relevancy ranking calculation for a multiple-word full-text search based on the `articles` table and data used in the previous example.

If you search on more than one word, the relevancy ranking value is a sum of the relevancy ranking value for each word, as shown in this formula:

```
${rank} = ${TF} * ${IDF} * ${IDF} + ${TF} * ${IDF} * ${IDF}
```

Performing a search on two terms ('mysql tutorial') returns the following results:

```
mysql> SELECT id, title, body, MATCH (title,body)  AGAINST ('mysql tutorial' IN BOOLEAN MODE)
    AS score FROM articles ORDER BY score DESC;
+----+-----------------------------+-----------------------------------+----------------------+
| id | title                       | body                              | score                |
+----+-----------------------------+-----------------------------------+----------------------+
|  1 | MySQL Tutorial              | This database tutorial ...        |   0.7405621409416199 |
|  3 | Optimizing Your Database    | In this database tutorial ...     |   0.3624762296676636 |
|  5 | MySQL Security              | When configured properly, MySQL ... | 0.031219376251101494 |
|  8 | MySQL Full-Text Indexes     | MySQL fulltext indexes use a ..   | 0.031219376251101494 |
|  2 | How To Use MySQL            | After you went through a ...      | 0.015609688125550747 |
|  4 | MySQL vs. YourSQL           | When comparing databases ...      | 0.015609688125550747 |
|  7 | 1001 MySQL Tricks           | 1. Never run mysqld as root. 2. ... | 0.015609688125550747 |
|  6 | Database, Database, Database | database database database        |                    0 |
+----+-----------------------------+-----------------------------------+----------------------+
8 rows in set (0.00 sec)
```

In the first record (`id 8`), 'mysql' appears once and 'tutorial' appears twice. There are six matching records for 'mysql' and two matching records for 'tutorial'. The MySQL command-line client returns the expected ranking value when inserting these values into the ranking formula for a multiple word search:

```
mysql> SELECT (1*log10(8/6)*log10(8/6)) + (2*log10(8/2)*log10(8/2));
+-----------------------------------------------------+
| (1*log10(8/6)*log10(8/6)) + (2*log10(8/2)*log10(8/2)) |
+-----------------------------------------------------+
|                                  0.7405621541938003 |
+-----------------------------------------------------+
1 row in set (0.00 sec)
```

Note

The slight difference in the ranking values returned by the `SELECT ... MATCH ... AGAINST` statement and the MySQL command-line client is explained in the preceding example.

12.9.3 Full-Text Searches with Query Expansion

Full-text search supports query expansion (and in particular, its variant “blind query expansion”). This is generally useful when a search phrase is too short, which often means that the user is relying on implied knowledge that the full-text search engine lacks. For example, a user searching for “database” may really mean that “MySQL”, “Oracle”, “DB2”, and “RDBMS” all are phrases that should match “databases” and should be returned, too. This is implied knowledge.

Blind query expansion (also known as automatic relevance feedback) is enabled by adding `WITH QUERY EXPANSION` or `IN NATURAL LANGUAGE MODE WITH QUERY EXPANSION` following the search phrase. It works by performing the search twice, where the search phrase for the second search is the original search phrase concatenated with the few most highly relevant documents from the first search. Thus, if one of these documents contains the word “databases” and the word “MySQL”, the second search finds the documents that contain the word “MySQL” even if they do not contain the word “database”. The following example shows this difference:

```
mysql> SELECT * FROM articles
    WHERE MATCH (title,body)
    AGAINST ('database' IN NATURAL LANGUAGE MODE);
+----+-------------------+----------------------------------------+
| id | title             | body                                   |
+----+-------------------+----------------------------------------+
|  1 | MySQL Tutorial    | DBMS stands for DataBase ...           |
|  5 | MySQL vs. YourSQL | In the following database comparison ... |
+----+-------------------+----------------------------------------+
2 rows in set (0.00 sec)

mysql> SELECT * FROM articles
    WHERE MATCH (title,body)
    AGAINST ('database' WITH QUERY EXPANSION);
+----+----------------------+----------------------------------------+
| id | title                | body                                   |
+----+----------------------+----------------------------------------+
|  5 | MySQL vs. YourSQL    | In the following database comparison ... |
|  1 | MySQL Tutorial       | DBMS stands for DataBase ...           |
|  3 | Optimizing MySQL     | In this tutorial we will show ...      |
|  6 | MySQL Security       | When configured properly, MySQL ...    |
|  2 | How To Use MySQL Well | After you went through a ...           |
|  4 | 1001 MySQL Tricks    | 1. Never run mysqld as root. 2. ...    |
+----+----------------------+----------------------------------------+
6 rows in set (0.00 sec)
```

Another example could be searching for books by Georges Simenon about Maigret, when a user is not sure how to spell “Maigret”. A search for “Megre and the reluctant witnesses” finds only “Maigret and the Reluctant Witnesses” without query expansion. A search with query expansion finds all books with the word “Maigret” on the second pass.

Note

Because blind query expansion tends to increase noise significantly by returning nonrelevant documents, use it only when a search phrase is short.

12.9.4 Full-Text Stopwords

The stopword list is loaded and searched for full-text queries using the server character set and collation (the values of the `character_set_server` and `collation_server` system variables). False hits or misses might occur for stopword lookups if the stopword file or columns used for full-text indexing or searches have a character set or collation different from `character_set_server` or `collation_server`.

Case sensitivity of stopword lookups depends on the server collation. For example, lookups are case insensitive if the collation is `utf8mb4_0900_ai_ci`, whereas lookups are case sensitive if the collation is `utf8mb4_0900_as_cs` or `utf8mb4_bin`.

Stopwords for InnoDB Search Indexes

InnoDB has a relatively short list of default stopwords, because documents from technical, literary, and other sources often use short words as keywords or in significant phrases. For example, you might search for “to be or not to be” and expect to get a sensible result, rather than having all those words ignored.

To see the default InnoDB stopword list, query the INFORMATION_SCHEMA.INNODB_FT_DEFAULT_STOPWORD table.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_DEFAULT_STOPWORD;
+-------+
| value |
+-------+
| a     |
| about |
| an    |
| are   |
| as    |
| at    |
| be    |
| by    |
| com   |
| de    |
| en    |
| for   |
| from  |
| how   |
| i     |
| in    |
| is    |
| it    |
| la    |
| of    |
| on    |
| or    |
| that  |
| the   |
| this  |
| to    |
| was   |
| what  |
| when  |
| where |
| who   |
| will  |
| with  |
| und   |
| the   |
| www   |
+-------+
36 rows in set (0.00 sec)
```

To define your own stopword list for all InnoDB tables, define a table with the same structure as the INNODB_FT_DEFAULT_STOPWORD table, populate it with stopwords, and set the value of the innodb_ft_server_stopword_table option to a value in the form *db_name/table_name* before creating the full-text index. The stopword table must have a single VARCHAR column named value. The following example demonstrates creating and configuring a new global stopword table for InnoDB.

```
-- Create a new stopword table

mysql> CREATE TABLE my_stopwords(value VARCHAR(30)) ENGINE = INNODB;
Query OK, 0 rows affected (0.01 sec)

-- Insert stopwords (for simplicity, a single stopword is used in this example)

mysql> INSERT INTO my_stopwords(value) VALUES ('Ishmael');
Query OK, 1 row affected (0.00 sec)

-- Create the table
```

```
mysql> CREATE TABLE opening_lines (
id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
opening_line TEXT(500),
author VARCHAR(200),
title VARCHAR(200)
) ENGINE=InnoDB;
Query OK, 0 rows affected (0.01 sec)

-- Insert data into the table

mysql> INSERT INTO opening_lines(opening_line,author,title) VALUES
('Call me Ishmael.','Herman Melville','Moby-Dick'),
('A screaming comes across the sky.','Thomas Pynchon','Gravity\'s Rainbow'),
('I am an invisible man.','Ralph Ellison','Invisible Man'),
('Where now? Who now? When now?','Samuel Beckett','The Unnamable'),
('It was love at first sight.','Joseph Heller','Catch-22'),
('All this happened, more or less.','Kurt Vonnegut','Slaughterhouse-Five'),
('Mrs. Dalloway said she would buy the flowers herself.','Virginia Woolf','Mrs. Dalloway'),
('It was a pleasure to burn.','Ray Bradbury','Fahrenheit 451');
Query OK, 8 rows affected (0.00 sec)
Records: 8  Duplicates: 0  Warnings: 0

-- Set the innodb_ft_server_stopword_table option to the new stopword table

mysql> SET GLOBAL innodb_ft_server_stopword_table = 'test/my_stopwords';
Query OK, 0 rows affected (0.00 sec)

-- Create the full-text index (which rebuilds the table if no FTS_DOC_ID column is defined)

mysql> CREATE FULLTEXT INDEX idx ON opening_lines(opening_line);
Query OK, 0 rows affected, 1 warning (1.17 sec)
Records: 0  Duplicates: 0  Warnings: 1
```

Verify that the specified stopword ('Ishmael') does not appear by querying the words in `INFORMATION_SCHEMA.INNODB_FT_INDEX_TABLE`.

Note

By default, words less than 3 characters in length or greater than 84 characters in length do not appear in an `InnoDB` full-text search index. Maximum and minimum word length values are configurable using the `innodb_ft_max_token_size` and `innodb_ft_min_token_size` variables. This default behavior does not apply to the ngram parser plugin. ngram token size is defined by the `ngram_token_size` option.

```
mysql> SET GLOBAL innodb_ft_aux_table='test/opening_lines';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT word FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_TABLE LIMIT 15;
+-----------+
| word      |
+-----------+
| across    |
| all       |
| burn      |
| buy       |
| call      |
| comes     |
| dalloway  |
| first     |
| flowers   |
| happened  |
| herself   |
| invisible |
| less      |
| love      |
| man       |
+-----------+
15 rows in set (0.00 sec)
```

To create stopword lists on a table-by-table basis, create other stopword tables and use the `innodb_ft_user_stopword_table` option to specify the stopword table that you want to use before you create the full-text index.

Stopwords for MyISAM Search Indexes

The stopword file is loaded and searched using `latin1` if `character_set_server` is `ucs2`, `utf16`, `utf16le`, or `utf32`.

To override the default stopword list for MyISAM tables, set the `ft_stopword_file` system variable. (See Section 5.1.5, “Server System Variables”.) The variable value should be the path name of the file containing the stopword list, or the empty string to disable stopword filtering. The server looks for the file in the data directory unless an absolute path name is given to specify a different directory. After changing the value of this variable or the contents of the stopword file, restart the server and rebuild your `FULLTEXT` indexes.

The stopword list is free-form, separating stopwords with any nonalphanumeric character such as newline, space, or comma. Exceptions are the underscore character (_) and a single apostrophe (') which are treated as part of a word. The character set of the stopword list is the server's default character set; see Section 10.1.3.2, “Server Character Set and Collation”.

The following table shows the default list of stopwords for `MyISAM` search indexes. In a MySQL source distribution, you can find this list in the `storage/myisam/ft_static.c` file.

a's	able	about	above	according
accordingly	across	actually	after	afterwards
again	against	ain't	all	allow
allows	almost	alone	along	already
also	although	always	am	among
amongst	an	and	another	any
anybody	anyhow	anyone	anything	anyway
anyways	anywhere	apart	appear	appreciate
appropriate	are	aren't	around	as
aside	ask	asking	associated	at
available	away	awfully	be	became
because	become	becomes	becoming	been
before	beforehand	behind	being	believe
below	beside	besides	best	better
between	beyond	both	brief	but
by	c'mon	c's	came	can
can't	cannot	cant	cause	causes
certain	certainly	changes	clearly	co
com	come	comes	concerning	consequently
consider	considering	contain	containing	contains
corresponding	could	couldn't	course	currently
definitely	described	despite	did	didn't
different	do	does	doesn't	doing
don't	done	down	downwards	during
each	edu	eg	eight	either
else	elsewhere	enough	entirely	especially

et	etc	even	ever	every
everybody	everyone	everything	everywhere	ex
exactly	example	except	far	few
fifth	first	five	followed	following
follows	for	former	formerly	forth
four	from	further	furthermore	get
gets	getting	given	gives	go
goes	going	gone	got	gotten
greetings	had	hadn't	happens	hardly
has	hasn't	have	haven't	having
he	he's	hello	help	hence
her	here	here's	hereafter	hereby
herein	hereupon	hers	herself	hi
him	himself	his	hither	hopefully
how	howbeit	however	i'd	i'll
i'm	i've	ie	if	ignored
immediate	in	inasmuch	inc	indeed
indicate	indicated	indicates	inner	insofar
instead	into	inward	is	isn't
it	it'd	it'll	it's	its
itself	just	keep	keeps	kept
know	known	knows	last	lately
later	latter	latterly	least	less
lest	let	let's	like	liked
likely	little	look	looking	looks
ltd	mainly	many	may	maybe
me	mean	meanwhile	merely	might
more	moreover	most	mostly	much
must	my	myself	name	namely
nd	near	nearly	necessary	need
needs	neither	never	nevertheless	new
next	nine	no	nobody	non
none	noone	nor	normally	not
nothing	novel	now	nowhere	obviously
of	off	often	oh	ok
okay	old	on	once	one
ones	only	onto	or	other
others	otherwise	ought	our	ours
ourselves	out	outside	over	overall
own	particular	particularly	per	perhaps
placed	please	plus	possible	presumably
probably	provides	que	quite	qv

rather	rd	re	really	reasonably
regarding	regardless	regards	relatively	respectively
right	said	same	saw	say
saying	says	second	secondly	see
seeing	seem	seemed	seeming	seems
seen	self	selves	sensible	sent
serious	seriously	seven	several	shall
she	should	shouldn't	since	six
so	some	somebody	somehow	someone
something	sometime	sometimes	somewhat	somewhere
soon	sorry	specified	specify	specifying
still	sub	such	sup	sure
t's	take	taken	tell	tends
th	than	thank	thanks	thanx
that	that's	thats	the	their
theirs	them	themselves	then	thence
there	there's	thereafter	thereby	therefore
therein	theres	thereupon	these	they
they'd	they'll	they're	they've	think
third	this	thorough	thoroughly	those
though	three	through	throughout	thru
thus	to	together	too	took
toward	towards	tried	tries	truly
try	trying	twice	two	un
under	unfortunately	unless	unlikely	until
unto	up	upon	us	use
used	useful	uses	using	usually
value	various	very	via	viz
vs	want	wants	was	wasn't
way	we	we'd	we'll	we're
we've	welcome	well	went	were
weren't	what	what's	whatever	when
whence	whenever	where	where's	whereafter
whereas	whereby	wherein	whereupon	wherever
whether	which	while	whither	who
who's	whoever	whole	whom	whose
why	will	willing	wish	with
within	without	won't	wonder	would
wouldn't	yes	yet	you	you'd
you'll	you're	you've	your	yours
yourself	yourselves	zero		

12.9.5 Full-Text Restrictions

- Full-text searches are supported for `InnoDB` and `MyISAM` tables only.
- Full-text searches are not supported for partitioned tables. See Section 22.6, “Restrictions and Limitations on Partitioning”.
- Full-text searches can be used with most multibyte character sets. The exception is that for Unicode, the `utf8` character set can be used, but not the `ucs2` character set. Although `FULLTEXT` indexes on `ucs2` columns cannot be used, you can perform `IN BOOLEAN MODE` searches on a `ucs2` column that has no such index.

  The remarks for `utf8` also apply to `utf8mb4`, and the remarks for `ucs2` also apply to `utf16`, `utf16le`, and `utf32`.
- Ideographic languages such as Chinese and Japanese do not have word delimiters. Therefore, the built-in full-text parser *cannot determine where words begin and end in these and other such languages*.

  A character-based ngram full-text parser that supports Chinese, Japanese, and Korean (CJK), and a word-based MeCab parser plugin that supports Japanese are provided for use with `InnoDB` and `MyISAM` tables.
- Although the use of multiple character sets within a single table is supported, all columns in a `FULLTEXT` index must use the same character set and collation.
- The `MATCH()` column list must match exactly the column list in some `FULLTEXT` index definition for the table, unless this `MATCH()` is `IN BOOLEAN MODE` on a `MyISAM` table. For `MyISAM` tables, boolean-mode searches can be done on nonindexed columns, although they are likely to be slow.
- The argument to `AGAINST()` must be a string value that is constant during query evaluation. This rules out, for example, a table column because that can differ for each row.
- Index hints are more limited for `FULLTEXT` searches than for non-`FULLTEXT` searches. See Section 8.9.4, “Index Hints”.
- For `InnoDB`, all DML operations (`INSERT`, `UPDATE`, `DELETE`) involving columns with full-text indexes are processed at transaction commit time. For example, for an `INSERT` operation, an inserted string is tokenized and decomposed into individual words. The individual words are then added to full-text index tables when the transaction is committed. As a result, full-text searches only return committed data.
- The '%' character is not a supported wildcard character for full-text searches.

12.9.6 Fine-Tuning MySQL Full-Text Search

MySQL's full-text search capability has few user-tunable parameters. You can exert more control over full-text searching behavior if you have a MySQL source distribution because some changes require source code modifications. See Section 2.8, “Installing MySQL from Source”.

Full-text search is carefully tuned for effectiveness. Modifying the default behavior in most cases can actually decrease effectiveness. *Do not alter the MySQL sources unless you know what you are doing*.

Most full-text variables described in this section must be set at server startup time. A server restart is required to change them; they cannot be modified while the server is running.

Some variable changes require that you rebuild the `FULLTEXT` indexes in your tables. Instructions for doing so are given later in this section.

- Configuring Minimum and Maximum Word Length

- Configuring the Natural Language Search Threshold
- Modifying Boolean Full-Text Search Operators
- Character Set Modifications
- Rebuilding InnoDB Full-Text Indexes
- Optimizing InnoDB Full-Text Indexes
- Rebuilding MyISAM Full-Text Indexes

Configuring Minimum and Maximum Word Length

The minimum and maximum lengths of words to be indexed are defined by the `innodb_ft_min_token_size` and `innodb_ft_max_token_size` for `InnoDB` search indexes, and `ft_min_word_len` and `ft_max_word_len` for `MyISAM` ones.

Note

Minimum and maximum word length full-text parameters do not apply to `FULLTEXT` indexes created using the ngram parser. ngram token size is defined by the `ngram_token_size` option.

After changing any of these options, rebuild your `FULLTEXT` indexes for the change to take effect. For example, to make two-character words searchable, you could put the following lines in an option file:

```
[mysqld]
innodb_ft_min_token_size=2
ft_min_word_len=2
```

Then restart the server and rebuild your `FULLTEXT` indexes. For `MyISAM` tables, note the remarks regarding `myisamchk` in the instructions that follow for rebuilding `MyISAM` full-text indexes.

Configuring the Natural Language Search Threshold

For `MyISAM` search indexes, the 50% threshold for natural language searches is determined by the particular weighting scheme chosen. To disable it, look for the following line in `storage/myisam/ftdefs.h`:

```
#define GWS_IN_USE GWS_PROB
```

Change that line to this:

```
#define GWS_IN_USE GWS_FREQ
```

Then recompile MySQL. There is no need to rebuild the indexes in this case.

Note

By making this change, you *severely* decrease MySQL's ability to provide adequate relevance values for the `MATCH()` function. If you really need to search for such common words, it would be better to search using `IN BOOLEAN MODE` instead, which does not observe the 50% threshold.

Modifying Boolean Full-Text Search Operators

To change the operators used for boolean full-text searches on `MyISAM` tables, set the `ft_boolean_syntax` system variable. (`InnoDB` does not have an equivalent setting.) This variable can be changed while the server is running, but you must have the `SYSTEM_VARIABLES_ADMIN` or

SUPER privilege to do so. No rebuilding of indexes is necessary in this case. See Section 5.1.5, “Server System Variables”, which describes the rules governing how to set this variable.

Character Set Modifications

For the built-in full-text parser, you can change the set of characters that are considered word characters in several ways, as described in the following list. After making the modification, rebuild the indexes for each table that contains any FULLTEXT indexes. Suppose that you want to treat the hyphen character ('-') as a word character. Use one of these methods:

- Modify the MySQL source: In storage/innobase/handler/ha_innodb.cc (for InnoDB), or in storage/myisam/ftdefs.h (for MyISAM), see the true_word_char() and misc_word_char() macros. Add '-' to one of those macros and recompile MySQL.
- Modify a character set file: This requires no recompilation. The true_word_char() macro uses a “character type” table to distinguish letters and numbers from other characters. . You can edit the contents of the <ctype><map> array in one of the character set XML files to specify that '-' is a “letter.” Then use the given character set for your FULLTEXT indexes. For information about the <ctype><map> array format, see Section 10.3.1, “Character Definition Arrays”.
- Add a new collation for the character set used by the indexed columns, and alter the columns to use that collation. For general information about adding collations, see Section 10.4, “Adding a Collation to a Character Set”. For an example specific to full-text indexing, see Section 12.9.7, “Adding a Collation for Full-Text Indexing”.

Rebuilding InnoDB Full-Text Indexes

If you modify full-text variables that affect indexing (innodb_ft_min_token_size, innodb_ft_max_token_size, innodb_ft_server_stopword_table, innodb_ft_user_stopword_table, innodb_ft_enable_stopword, ngram_token_size you must rebuild your FULLTEXT indexes after making the changes. Modifying the innodb_ft_min_token_size, innodb_ft_max_token_size, or ngram_token_size variables, which cannot be set dynamically, require restarting the server and rebuilding the indexes.

To rebuild the FULLTEXT indexes for an InnoDB table, use ALTER TABLE with the DROP INDEX and ADD INDEX options to drop and re-create each index.

Optimizing InnoDB Full-Text Indexes

Running OPTIMIZE TABLE on a table with a full-text index rebuilds the full-text index, removing deleted Document IDs and consolidating multiple entries for the same word, where possible.

To optimize a full-text index, enable innodb_optimize_fulltext_only and run OPTIMIZE TABLE.

```
mysql> set GLOBAL innodb_optimize_fulltext_only=ON;
Query OK, 0 rows affected (0.01 sec)

mysql> OPTIMIZE TABLE opening_lines;
+--------------------+----------+----------+----------+
| Table              | Op       | Msg_type | Msg_text |
+--------------------+----------+----------+----------+
| test.opening_lines | optimize | status   | OK       |
+--------------------+----------+----------+----------+
1 row in set (0.01 sec)
```

To avoid lengthy rebuild times for full-text indexes on large tables, you can use the innodb_ft_num_word_optimize option to perform the optimization in stages. The innodb_ft_num_word_optimize option defines the number of words that are optimized each time OPTIMIZE TABLE is run. The default setting is 2000, which means that 2000 words are optimized each time OPTIMIZE TABLE is run. Subsequent OPTIMIZE TABLE operations continue from where the preceding OPTIMIZE TABLE operation ended.

Rebuilding MyISAM Full-Text Indexes

If you modify full-text variables that affect indexing (`ft_min_word_len`, `ft_max_word_len`, or `ft_stopword_file`), or if you change the stopword file itself, you must rebuild your `FULLTEXT` indexes after making the changes and restarting the server.

To rebuild the `FULLTEXT` indexes for a `MyISAM` table, it is sufficient to do a `QUICK` repair operation:

```
mysql> REPAIR TABLE tbl_name QUICK;
```

Alternatively, use `ALTER TABLE` as just described. In some cases, this may be faster than a repair operation.

Each table that contains any `FULLTEXT` index must be repaired as just shown. Otherwise, queries for the table may yield incorrect results, and modifications to the table will cause the server to see the table as corrupt and in need of repair.

If you use `myisamchk` to perform an operation that modifies `MyISAM` table indexes (such as repair or analyze), the `FULLTEXT` indexes are rebuilt using the *default* full-text parameter values for minimum word length, maximum word length, and stopword file unless you specify otherwise. This can result in queries failing.

The problem occurs because these parameters are known only by the server. They are not stored in `MyISAM` index files. To avoid the problem if you have modified the minimum or maximum word length or stopword file values used by the server, specify the same `ft_min_word_len`, `ft_max_word_len`, and `ft_stopword_file` values for `myisamchk` that you use for `mysqld`. For example, if you have set the minimum word length to 3, you can repair a table with `myisamchk` like this:

```
shell> myisamchk --recover --ft_min_word_len=3 tbl_name.MYI
```

To ensure that `myisamchk` and the server use the same values for full-text parameters, place each one in both the `[mysqld]` and `[myisamchk]` sections of an option file:

```
[mysqld]
ft_min_word_len=3

[myisamchk]
ft_min_word_len=3
```

An alternative to using `myisamchk` for `MyISAM` table index modification is to use the `REPAIR TABLE`, `ANALYZE TABLE`, `OPTIMIZE TABLE`, or `ALTER TABLE` statements. These statements are performed by the server, which knows the proper full-text parameter values to use.

12.9.7 Adding a Collation for Full-Text Indexing

This section describes how to add a new collation for full-text searches using the built-in full-text parser. The sample collation is like `latin1_swedish_ci` but treats the `'-'` character as a letter rather than as a punctuation character so that it can be indexed as a word character. General information about adding collations is given in Section 10.4, “Adding a Collation to a Character Set”; it is assumed that you have read it and are familiar with the files involved.

To add a collation for full-text indexing, use the following procedure. The instructions here add a collation for a simple character set, which as discussed in Section 10.4, “Adding a Collation to a Character Set”, can be created using a configuration file that describes the character set properties. For a complex character set such as Unicode, create collations using C source files that describe the character set properties.

1. Add a collation to the `Index.xml` file. The collation ID must be unused, so choose a value different from 1000 if that ID is already taken on your system.

```
<charset name="latin1">
...
<collation name="latin1_fulltext_ci" id="1000"/>
</charset>
```

2. Declare the sort order for the collation in the `latin1.xml` file. In this case, the order can be copied from `latin1_swedish_ci`:

```
<collation name="latin1_fulltext_ci">
<map>
00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F
20 21 22 23 24 25 26 27 28 29 2A 2B 2C 2D 2E 2F
30 31 32 33 34 35 36 37 38 39 3A 3B 3C 3D 3E 3F
40 41 42 43 44 45 46 47 48 49 4A 4B 4C 4D 4E 4F
50 51 52 53 54 55 56 57 58 59 5A 5B 5C 5D 5E 5F
60 41 42 43 44 45 46 47 48 49 4A 4B 4C 4D 4E 4F
50 51 52 53 54 55 56 57 58 59 5A 7B 7C 7D 7E 7F
80 81 82 83 84 85 86 87 88 89 8A 8B 8C 8D 8E 8F
90 91 92 93 94 95 96 97 98 99 9A 9B 9C 9D 9E 9F
A0 A1 A2 A3 A4 A5 A6 A7 A8 A9 AA AB AC AD AE AF
B0 B1 B2 B3 B4 B5 B6 B7 B8 B9 BA BB BC BD BE BF
41 41 41 41 5C 5B 5C 43 45 45 45 45 49 49 49 49
44 4E 4F 4F 4F 4F 5D D7 D8 55 55 55 59 59 DE DF
41 41 41 41 5C 5B 5C 43 45 45 45 45 49 49 49 49
44 4E 4F 4F 4F 4F 5D F7 D8 55 55 55 59 59 DE FF
</map>
</collation>
```

3. Modify the `ctype` array in `latin1.xml`. Change the value corresponding to 0x2D (which is the code for the `'-'` character) from 10 (punctuation) to 01 (small letter). In the following array, this is the element in the fourth row down, third value from the end.

```
<ctype>
<map>
00
20 20 20 20 20 20 20 20 20 28 28 28 28 28 20 20
20 20 20 20 20 20 20 20 20 20 20 20 20 20 20 20
48 10 10 10 10 10 10 10 10 10 10 10 01 10 10 10
84 84 84 84 84 84 84 84 84 10 10 10 10 10 10 10
10 81 81 81 81 81 81 01 01 01 01 01 01 01 01 01
01 01 01 01 01 01 01 01 01 01 10 10 10 10 10 10
10 82 82 82 82 82 82 02 02 02 02 02 02 02 02 02
02 02 02 02 02 02 02 02 02 02 10 10 10 10 20 20
10 00 10 02 10 10 10 10 10 10 01 10 01 00 01 00
00 10 10 10 10 10 10 10 10 10 02 10 02 00 02 01
48 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10
10 10 10 10 10 10 10 10 10 10 10 10 10 10 10 10
01 01 01 01 01 01 01 01 01 01 01 01 01 01 01 01
01 01 01 01 01 01 01 10 01 01 01 01 01 01 01 02
02 02 02 02 02 02 02 02 02 02 02 02 02 02 02 02
02 02 02 02 02 02 02 10 02 02 02 02 02 02 02 02
</map>
</ctype>
```

4. Restart the server.

5. To employ the new collation, include it in the definition of columns that are to use it:

```
mysql> DROP TABLE IF EXISTS t1;
Query OK, 0 rows affected (0.13 sec)

mysql> CREATE TABLE t1 (
    a TEXT CHARACTER SET latin1 COLLATE latin1_fulltext_ci,
    FULLTEXT INDEX(a)
    ) ENGINE=InnoDB;
```

```
Query OK, 0 rows affected (0.47 sec)
```

6. Test the collation to verify that hyphen is considered as a word character:

```
mysql> INSERT INTO t1 VALUEs ('----'),('....'),('abcd');
Query OK, 3 rows affected (0.22 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM t1 WHERE MATCH a AGAINST ('----' IN BOOLEAN MODE);
+------+
| a    |
+------+
| ---- |
+------+
1 row in set (0.00 sec)
```

12.9.8 ngram Full-Text Parser

The built-in MySQL full-text parser uses the white space between words as a delimiter to determine where words begin and end, which is a limitation when working with ideographic languages that do not use word delimiters. To address this limitation, MySQL provides an ngram full-text parser that supports Chinese, Japanese, and Korean (CJK). The ngram full-text parser is supported for use with `InnoDB` and `MyISAM`.

Note

MySQL also provides a MeCab full-text parser plugin for Japanese, which tokenizes documents into meaningful words. For more information, see Section 12.9.9, “MeCab Full-Text Parser Plugin”.

An ngram is a contiguous sequence of *n* characters from a given sequence of text. The ngram parser tokenizes a sequence of text into a contiguous sequence of *n* characters. For example, you can tokenize “abcd” for different values of *n* using the ngram full-text parser.

```
n=1: 'a', 'b', 'c', 'd'
n=2: 'ab', 'bc', 'cd'
n=3: 'abc', 'bcd'
n=4: 'abcd'
```

The ngram full-text parser is a built-in server plugin. As with other built-in server plugins, it is automatically loaded when the server is started.

The full-text search syntax described in Section 12.9, “Full-Text Search Functions” applies to the ngram parser plugin. Differences in parsing behavior are described in this section. Full-text-related configuration options, except for minimum and maximum word length options (`innodb_ft_min_token_size`, `innodb_ft_max_token_size`, `ft_min_word_len`, `ft_max_word_len`) are also applicable.

Configuring ngram Token Size

The ngram parser has a default ngram token size of 2 (bigram). For example, with a token size of 2, the ngram parser parses the string “abc def” into four tokens: “ab”, “bc”, “de” and “ef”.

ngram token size is configurable using the `ngram_token_size` configuration option, which has a minimum value of 1 and maximum value of 10.

Typically, `ngram_token_size` is set to the size of the largest token that you want to search for. If you only intend to search for single characters, set `ngram_token_size` to 1. A smaller token size produces a smaller full-text search index, and faster searches. If you need to search for words comprised of more than one character, set `ngram_token_size` accordingly. For example, “Happy Birthday” is “生日快乐” in simplified Chinese, where “生日” is “birthday”, and “快乐” translates as “happy”. To search on two-character words such as these, set `ngram_token_size` to a value of 2 or higher.

As a read-only variable, `ngram_token_size` may only be set as part of a startup string or in a configuration file:

- Startup string:

```
mysqld --ngram_token_size=2
```

- Configuration file:

```
[mysqld]
ngram_token_size=2
```

Note

The following minimum and maximum word length configuration options are ignored for `FULLTEXT` indexes that use the ngram parser: `innodb_ft_min_token_size`, `innodb_ft_max_token_size`, `ft_min_word_len`, and `ft_max_word_len`.

Creating a FULLTEXT Index that Uses the ngram Parser

To create a `FULLTEXT` index that uses the ngram parser, specify `WITH PARSER ngram` with `CREATE TABLE`, `ALTER TABLE`, or `CREATE INDEX`.

The following example demonstrates creating a table with an `ngram` `FULLTEXT` index, inserting sample data (Simplified Chinese text), and viewing tokenized data in the `INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE` table.

```
mysql> USE test;

mysql> CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT,
      FULLTEXT (title,body) WITH PARSER ngram
    ) ENGINE=InnoDB CHARACTER SET utf8mb4;

mysql> SET NAMES utf8mb4;

INSERT INTO articles (title,body) VALUES
    ('数据库管理','在本教程中我将向你展示如何管理数据库'),
    ('数据库应用开发','学习开发数据库应用程序');

mysql> SET GLOBAL innodb_ft_aux_table="test/articles";

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE ORDER BY doc_id, position;
```

To add a `FULLTEXT` index to an existing table, you can use `ALTER TABLE` or `CREATE INDEX`. For example:

```
CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT
     ) ENGINE=InnoDB CHARACTER SET utf8;

ALTER TABLE articles ADD FULLTEXT INDEX ft_index (title,body) WITH PARSER ngram;

# Or:

CREATE FULLTEXT INDEX ft_index ON articles (title,body) WITH PARSER ngram;
```

ngram Parser Space Handling

The ngram parser eliminates spaces when parsing. For example:

- “ab cd” is parsed to “ab”, “cd”

- “a bc” is parsed to “bc”

ngram Parser Stopword Handling

The built-in MySQL full-text parser compares words to entries in the stopword list. If a word is equal to an entry in the stopword list, the word is excluded from the index. For the ngram parser, stopword handling is performed differently. Instead of excluding tokens that are equal to entries in the stopword list, the ngram parser excludes tokens that *contain* stopwords. For example, assuming `ngram_token_size=2`, a document that contains “a,b” is parsed to “a,” and “,b”. If a comma (“,”) is defined as a stopword, both “a,” and “,b” are excluded from the index because they contain a comma.

By default, the ngram parser uses the default stopword list, which contains a list of English stopwords. For a stopword list applicable to Chinese, Japanese, or Korean, you must create your own. For information about creating a stopword list, see Section 12.9.4, “Full-Text Stopwords”.

Stopwords greater in length than `ngram_token_size` are ignored.

ngram Parser Term Search

For *natural language mode* search, the search term is converted to a union of ngram terms. For example, the string “abc” (assuming `ngram_token_size=2`) is converted to “ab bc”. Given two documents, one containing “ab” and the other containing “abc”, the search term “ab bc” matches both documents.

For *boolean mode search*, the search term is converted to an ngram phrase search. For example, the string 'abc' (assuming `ngram_token_size=2`) is converted to '“ab bc”'. Given two documents, one containing 'ab' and the other containing 'abc', the search phrase '“ab bc”' only matches the document containing 'abc'.

ngram Parser Wildcard Search

Because an ngram `FULLTEXT` index contains only ngrams, and does not contain information about the beginning of terms, wildcard searches may return unexpected results. The following behaviors apply to wildcard searches using ngram `FULLTEXT` search indexes:

- If the prefix term of a wildcard search is shorter than ngram token size, the query returns all indexed rows that contain ngram tokens starting with the prefix term. For example, assuming `ngram_token_size=2`, a search on “a*” returns all rows starting with “a”.
- If the prefix term of a wildcard search is longer than ngram token size, the prefix term is converted to an ngram phrase and the wildcard operator is ignored. For example, assuming `ngram_token_size=2`, an “abc*” wildcard search is converted to “ab bc”.

ngram Parser Phrase Search

Phrase searches are converted to ngram phrase searches. For example, The search phrase “abc” is converted to “ab bc”, which returns documents containing “abc” and “ab bc”.

The search phrase “abc def” is converted to “ab bc de ef”, which returns documents containing “abc def” and “ab bc de ef”. A document that contains “abcdef” is not returned.

12.9.9 MeCab Full-Text Parser Plugin

The built-in MySQL full-text parser uses the white space between words as a delimiter to determine where words begin and end, which is a limitation when working with ideographic languages that do not use word delimiters. To address this limitation for Japanese, MySQL provides a MeCab full-text parser plugin. The MeCab full-text parser plugin is supported for use with `InnoDB` and `MyISAM`.

Note

MySQL also provides an ngram full-text parser plugin that supports Japanese. For more information, see Section 12.9.8, “ngram Full-Text Parser”.

The MeCab full-text parser plugin is a full-text parser plugin for Japanese that tokenizes a sequence of text into meaningful words. For example, MeCab tokenizes “データベース管理” (“Database Management”) into “データベース” (“Database”) and “管理” (“Management”). By comparison, the ngram full-text parser tokenizes text into a contiguous sequence of *n* characters, where *n* represents a number between 1 and 10.

In addition to tokenizing text into meaningful words, MeCab indexes are typically smaller than ngram indexes, and MeCab full-text searches are generally faster. One drawback is that it may take longer for the MeCab full-text parser to tokenize documents, compared to the ngram full-text parser.

The full-text search syntax described in Section 12.9, “Full-Text Search Functions” applies to the MeCab parser plugin. Differences in parsing behavior are described in this section. Full-text related configuration options are also applicable.

For additional information about the MeCab parser, refer to the MeCab Documentation on the Google Developers site.

Installing the MeCab Parser Plugin

The MeCab parser plugin requires `mecab` and `mecab-ipadic`.

On supported Fedora, Debian and Ubuntu platforms (except Ubuntu 12.04 where the system `mecab` version is too old), MySQL dynamically links to the system `mecab` installation if it is installed to the default location. On other supported Unix-like platforms, `libmecab.so` is statically linked in `libpluginmecab.so`, which is located in the MySQL plugin directory. `mecab-ipadic` is included in MySQL binaries and is located in *`MYSQL_HOME`*`\lib\mecab`.

You can install `mecab` and `mecab-ipadic` using a native package management utility (on Fedora, Debian, and Ubuntu), or you can build `mecab` and `mecab-ipadic` from source. For information about installing `mecab` and `mecab-ipadic` using a native package management utility, see Installing MeCab From a Binary Distribution (Optional). If you want to build `mecab` and `mecab-ipadic` from source, see Building MeCab From Source (Optional).

On Windows, `libmecab.dll` is found in the MySQL `bin` directory. `mecab-ipadic` is located in *`MYSQL_HOME`*`/lib/mecab`.

To install and configure the MeCab parser plugin, perform the following steps:

1. In the MySQL configuration file, set the `mecab_rc_file` configuration option to the location of the `mecabrc` configuration file, which is the configuration file for MeCab. If you are using the MeCab package distributed with MySQL, the `mecabrc` file is located in `MYSQL_HOME/lib/mecab/etc/`.

   ```
   [mysqld]
   loose-mecab-rc-file=MYSQL_HOME/lib/mecab/etc/mecabrc
   ```

   The `loose` prefix is an option modifier. The `mecab_rc_file` option is not recognized by MySQL until the MeCaB parser plugin is installed but it must be set before attempting to install the MeCaB parser plugin. The `loose` prefix allows you restart MySQL without encountering an error due to an unrecognized variable.

   If you use your own MeCab installation, or build MeCab from source, the location of the `mecabrc` configuration file may differ.

   For information about the MySQL configuration file and its location, see Section 4.2.6, “Using Option Files”.

2. Also in the MySQL configuration file, set the minimum token size to 1 or 2, which are the values recommended for use with the MeCab parser. For `InnoDB` tables, minimum token size is defined by the `innodb_ft_min_token_size` configuration option, which has a default value of 3. For `MyISAM` tables, minimum token size is defined by `ft_min_word_len`, which has a default value of 4.

```
[mysqld]
innodb_ft_min_token_size=1
```

3. Modify the `mecabrc` configuration file to specify the dictionary you want to use. The `mecab-ipadic` package distributed with MySQL binaries includes three dictionaries (`ipadic_euc-jp`, `ipadic_sjis`, and `ipadic_utf-8`). The `mecabrc` configuration file packaged with MySQL contains and entry similar to the following:

   ```
   dicdir =  /path/to/mysql/lib/mecab/lib/mecab/dic/ipadic_euc-jp
   ```

   To use the `ipadic_utf-8` dictionary, for example, modify the entry as follows:

   ```
   dicdir=MYSQL_HOME/lib/mecab/dic/ipadic_utf-8
   ```

   If you are using your own MeCab installation or have built MeCab from source, the default `dicdir` entry in the `mecabrc` file will differ, as will the dictionaries and their location.

   Note

   After the MeCab parser plugin is installed, you can use the `mecab_charset` status variable to view the character set used with MeCab. The three MeCab dictionaries provided with the MySQL binary support the following character sets.

   - The `ipadic_euc-jp` dictionary supports the `ujis` and `eucjpms` character sets.
   - The `ipadic_sjis` dictionary supports the `sjis` and `cp932` character sets.
   - The `ipadic_utf-8` dictionary supports the `utf8` and `utf8mb4` character sets.

   `mecab_charset` only reports the first supported character set. For example, the `ipadic_utf-8` dictionary supports both `utf8` and `utf8mb4`. `mecab_charset` always reports `utf8` when this dictionary is in use.

4. Restart MySQL.

5. Install the MeCab parser plugin:

   The MeCab parser plugin is installed using `INSTALL PLUGIN` syntax. The plugin name is `mecab`, and the shared library name is `libpluginmecab.so`. For additional information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

   ```
   INSTALL PLUGIN mecab SONAME 'libpluginmecab.so';
   ```

   Once installed, the MeCab parser plugin loads at every normal MySQL restart.

6. Verify that the MeCab parser plugin is loaded using the `SHOW PLUGINS` statement.

   ```
   mysql> SHOW PLUGINS;
   ```

   A `mecab` plugin should appear in the list of plugins.

Creating a FULLTEXT Index that uses the MeCab Parser

To create a `FULLTEXT` index that uses the mecab parser, specify `WITH PARSER ngram` with `CREATE TABLE`, `ALTER TABLE`, or `CREATE INDEX`.

This example demonstrates creating a table with a `mecab` `FULLTEXT` index, inserting sample data, and viewing tokenized data in the `INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE` table:

```
mysql> USE test;

mysql> CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT,
      FULLTEXT (title,body) WITH PARSER mecab
    ) ENGINE=InnoDB CHARACTER SET utf8;

mysql> SET NAMES utf8;

mysql> INSERT INTO articles (title,body) VALUES
    ('データベース管理','このチュートリアルでは、私はどのようにデータベースを管理する方法を紹介します'),
    ('データベースアプリケーション開発','データベースアプリケーションを開発することを学ぶ');

mysql> SET GLOBAL innodb_ft_aux_table="test/articles";

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE ORDER BY doc_id, position;
```

To add a `FULLTEXT` index to an existing table, you can use `ALTER TABLE` or `CREATE INDEX`. For example:

```
CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT
     ) ENGINE=InnoDB CHARACTER SET utf8;

ALTER TABLE articles ADD FULLTEXT INDEX ft_index (title,body) WITH PARSER mecab;

# Or:

CREATE FULLTEXT INDEX ft_index ON articles (title,body) WITH PARSER mecab;
```

MeCab Parser Space Handling

The MeCab parser uses spaces as separators in query strings. For example, the MeCab parser tokenizes 'データベース管理' as 'データベース' and '管理'.

MeCab Parser Stopword Handling

By default, the MeCab parser uses the default stopword list, which contains a short list of English stopwords. For a stopword list applicable to Japanese, you must create your own. For information about creating stopword lists, see Section 12.9.4, “Full-Text Stopwords”.

MeCab Parser Term Search

For natural language mode search, the search term is converted to a union of tokens. For example, 'データベース管理' is converted to 'データベース 管理' .

```
SELECT COUNT(*) FROM articles WHERE MATCH(title,body) AGAINST('データベース管理' IN NATURAL LANGUAGE MODE);
```

For boolean mode search, the search term is converted to a search phrase. For example, 'データベース管理' is converted to '"データベース 管理"'.

```
SELECT COUNT(*) FROM articles WHERE MATCH(title,body) AGAINST('データベース管理' IN BOOLEAN MODE);
```

MeCab Parser Wildcard Search

Wildcard search terms are not tokenized. A search on 'データベース管理*' is performed on the prefix, 'データベース管理'.

```
SELECT COUNT(*) FROM articles WHERE MATCH(title,body) AGAINST('データベース*' IN BOOLEAN MODE);
```

MeCab Parser Phrase Search

Phrases are tokenized. For example, "データベース管理" is tokenized as "データベース 管理".

```
SELECT COUNT(*) FROM articles WHERE MATCH(title,body) AGAINST('"データベース管理"' IN BOOLEAN MODE);
```

Installing MeCab From a Binary Distribution (Optional)

This section describes how to install `mecab` and `mecab-ipadic` from a binary distribution using a native package management utility. For example, on Fedora, you can use Yum to perform the installation:

```
yum mecab-devel
```

On Debian or Ubuntu, you can perform an APT installation:

```
apt-get install mecab
apt-get install mecab-ipadic
```

Installing MeCab From Source (Optional)

If you want to build `mecab` and `mecab-ipadic` from source, basic installation steps are provided below. For additional information, refer to the MeCab documentation.

1. Download the tar.gz packages for `mecab` and `mecab-ipadic` from http://taku910.github.io/mecab/#download. As of February, 2016, the latest available packages are `mecab-0.996.tar.gz` and `mecab-ipadic-2.7.0-20070801.tar.gz`.

2. Install `mecab`:

   ```
   tar zxfv mecab-0.996.tar
   cd mecab-0.996
   ./configure
   make
   make check
   su
   make install
   ```

3. Install `mecab-ipadic`:

   ```
   tar zxfv mecab-ipadic-2.7.0-20070801.tar
   cd mecab-ipadic-2.7.0-20070801
   ./configure
   make
   su
   make install
   ```

4. Compile MySQL using the `WITH_MECAB` CMake option. Set the `WITH_MECAB` option to `system` if you have installed `mecab` and `mecab-ipadic` to the default location.

   ```
   -DWITH_MECAB=system
   ```

   If you defined a custom installation directory, set `WITH_MECAB` to the custom directory. For example:

   ```
   -DWITH_MECAB=/path/to/mecab
   ```

12.10 Cast Functions and Operators

Table 12.14 Cast Functions and Operators

Name	Description
`BINARY`	Cast a string to a binary string
`CAST()`	Cast a value as a certain type
`CONVERT()`	Cast a value as a certain type

Cast functions and operators enable conversion of values from one data type to another.

`CONVERT()` with a `USING` clause provides a way to convert data between different character sets:

```
CONVERT(expr USING transcoding_name)
```

In MySQL, transcoding names are the same as the corresponding character set names.

Examples:

```
SELECT CONVERT(_latin1'Müller' USING utf8);
INSERT INTO utf8_table (utf8_column)
    SELECT CONVERT(latin1_column USING utf8) FROM latin1_table;
```

You can also use `CONVERT()` without `USING` or `CAST()` to convert strings between different character sets:

```
CONVERT(string, CHAR[(N)] CHARACTER SET charset_name)
CAST(string AS CHAR[(N)] CHARACTER SET charset_name)
```

Examples:

```
SELECT CONVERT('test', CHAR CHARACTER SET utf8);
SELECT CAST('test' AS CHAR CHARACTER SET utf8);
```

If you specify `CHARACTER SET charset_name` as just shown, the resulting character set and collation are `charset_name` and the default collation of `charset_name`. If you omit `CHARACTER SET charset_name`, the resulting character set and collation are defined by the `character_set_connection` and `collation_connection` system variables that determine the default connection character set and collation (see Section 10.1.4, “Connection Character Sets and Collations”).

A `COLLATE` clause is not permitted within a `CONVERT()` or `CAST()` call, but you can apply it to the function result. For example, this is legal:

```
SELECT CAST('test' AS CHAR CHARACTER SET utf8) COLLATE utf8_bin;
```

But this is illegal:

```
SELECT CAST('test' AS CHAR CHARACTER SET utf8 COLLATE utf8_bin);
```

Normally, you cannot compare a `BLOB` value or other binary string in case-insensitive fashion because binary strings use the `binary` character set, which has no collation with the concept of lettercase. To perform a case-insensitive comparison, use the `CONVERT()` or `CAST()` function to convert the value to a nonbinary string. Comparisons of the resulting string use its collation. For example, if the conversion result character set has a case-insensitive collation, a `LIKE` operation is not case sensitive:

```
SELECT 'A' LIKE CONVERT(blob_col USING latin1)
  FROM tbl_name;
```

To use a different character set, substitute its name for `latin1` in the preceding statement. To specify a particular collation for the converted string, use a `COLLATE` clause following the `CONVERT()` call:

```
SELECT 'A' LIKE CONVERT(blob_col USING latin1) COLLATE latin1_german1_ci
  FROM tbl_name;
```

`CONVERT()` and `CAST()` can be used more generally for comparing strings that are represented in different character sets. For example, a comparison of these strings results in an error because they have different character sets:

```
```

```
mysql> SET @s1 = _latin1 'abc', @s2 = _latin2 'abc';
mysql> SELECT @s1 = @s2;
ERROR 1267 (HY000): Illegal mix of collations (latin1_swedish_ci,IMPLICIT)
and (latin2_general_ci,IMPLICIT) for operation '='
```

Converting one of the strings to a character set compatible with the other enables the comparison to occur without error:

```
mysql> SELECT @s1 = CONVERT(@s2 USING latin1);
+--------------------------------+
| @s1 = CONVERT(@s2 USING latin1) |
+--------------------------------+
|                              1 |
+--------------------------------+
```

For string literals, another way to specify the character set is to use a character set introducer (`_latin1` and `_latin2` in the preceding example are instances of introducers). Unlike conversion functions such as `CAST()`, or `CONVERT()`, which convert a string from one character set to another, an introducer designates a string literal as having a particular character set, with no conversion involved. For more information, see Section 10.1.3.8, “Character Set Introducers”.

Character set conversion is also useful preceding lettercase conversion of binary strings. `LOWER()` and `UPPER()` are ineffective when applied directly to binary strings because the concept of lettercase does not apply. To perform lettercase conversion of a binary string, first convert it to a nonbinary string:

```
mysql> SET @str = BINARY 'New York';
mysql> SELECT LOWER(@str), LOWER(CONVERT(@str USING utf8mb4));
+-------------+-----------------------------------+
| LOWER(@str) | LOWER(CONVERT(@str USING utf8mb4)) |
+-------------+-----------------------------------+
| New York    | new york                          |
+-------------+-----------------------------------+
```

If you convert an indexed column using `BINARY`, `CAST()`, or `CONVERT()`, MySQL may not be able to use the index efficiently.

The cast functions are useful for creating a column with a specific type in a `CREATE TABLE ... SELECT` statement:

```
mysql> CREATE TABLE new_table SELECT CAST('2000-01-01' AS DATE) AS c1;
mysql> SHOW CREATE TABLE new_table\G
*************************** 1. row ***************************
       Table: new_table
Create Table: CREATE TABLE `new_table` (
  `c1` date DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4
```

The cast functions are useful for sorting `ENUM` columns in lexical order. Normally, sorting of `ENUM` columns occurs using the internal numeric values. Casting the values to `CHAR` results in a lexical sort:

```
SELECT enum_col FROM tbl_name ORDER BY CAST(enum_col AS CHAR);
```

`CAST()` also changes the result if you use it as part of a more complex expression such as `CONCAT('Date: ',CAST(NOW() AS DATE))`.

For temporal values, there is little need to use `CAST()` to extract data in different formats. Instead, use a function such as `EXTRACT()`, `DATE_FORMAT()`, or `TIME_FORMAT()`. See Section 12.7, “Date and Time Functions”.

To cast a string to a number, you normally need do nothing other than use the string value in numeric context:

```
mysql> SELECT 1+'1';
       -> 2
```

That is also true for hexadecimal and bit literals, which are binary strings by default:

```
mysql> SELECT X'41', X'41'+0;
        -> 'A', 65
mysql> SELECT b'1100001', b'1100001'+0;
        -> 'a', 97
```

A string used in an arithmetic operation is converted to a floating-point number during expression evaluation.

A number used in string context is converted to a string:

```
mysql> SELECT CONCAT('hello you ',2);
        -> 'hello you 2'
```

For information about implicit conversion of numbers to strings, see Section 12.2, “Type Conversion in Expression Evaluation”.

MySQL supports arithmetic with both signed and unsigned 64-bit values. For numeric operators (such as `+` or `-`) where one of the operands is an unsigned integer, the result is unsigned by default (see Section 12.6.1, “Arithmetic Operators”). To override this, use the `SIGNED` or `UNSIGNED` cast operator to cast a value to a signed or unsigned 64-bit integer, respectively.

```
mysql> SELECT 1 - 2;
        -> -1
mysql> SELECT CAST(1 - 2 AS UNSIGNED);
        -> 18446744073709551615
mysql> SELECT CAST(CAST(1 - 2 AS UNSIGNED) AS SIGNED);
        -> -1
```

If either operand is a floating-point value, the result is a floating-point value and is not affected by the preceding rule. (In this context, `DECIMAL` column values are regarded as floating-point values.)

```
mysql> SELECT CAST(1 AS UNSIGNED) - 2.0;
        -> -1.0
```

The SQL mode affects the result of conversion operations (see Section 5.1.8, “Server SQL Modes”). Examples:

- For conversion of a “zero” date string to a date, `CONVERT()` and `CAST()` return `NULL` and produce a warning when the `NO_ZERO_DATE` SQL mode is enabled.
- For integer subtraction, if the `NO_UNSIGNED_SUBTRACTION` SQL mode is enabled, the subtraction result is signed even if any operand is unsigned.

The following list describes the available cast functions and operators:

- `BINARY` *expr*

  The `BINARY` operator converts the expression to a binary string. A common use for `BINARY` is to force a character string comparison to be done byte by byte rather than character by character, in effect becoming case sensitive. The `BINARY` operator also causes trailing spaces in comparisons to be significant.

  ```
  mysql> SELECT 'a' = 'A';
          -> 1
  mysql> SELECT BINARY 'a' = 'A';
          -> 0
  mysql> SELECT 'a' = 'a ';
  ```

```
        -> 1
mysql> SELECT BINARY 'a' = 'a ';
        -> 0
```

In a comparison, `BINARY` affects the entire operation; it can be given before either operand with the same result.

For purposes of converting a string expression to a binary string, these constructs are equivalent:

```
BINARY expr
CAST(expr AS BINARY)
CONVERT(expr USING BINARY)
```

If a value is a string literal, it can be designated as a binary string without performing any conversion by using the `_binary` character set introducer:

```
mysql> SELECT 'a' = 'A';
        -> 1
mysql> SELECT _binary 'a' = 'A';
        -> 0
```

For information about introducers, see Section 10.1.3.8, “Character Set Introducers”.

The `BINARY` operator in expressions differs in effect from the `BINARY` attribute in character column definitions. A character column defined with the `BINARY` attribute is assigned table default character set and the binary (`_bin`) collation of that character set. Every nonbinary character set has a `_bin` collation. For example, the binary collation for the `utf8` character set is `utf8_bin`, so if the table default character set is `utf8`, these two column definitions are equivalent:

```
CHAR(10) BINARY
CHAR(10) CHARACTER SET utf8 COLLATE utf8_bin
```

The use of `CHARACTER SET binary` in the definition of a `CHAR`, `VARCHAR`, or `TEXT` column causes the column to be treated as the corresponding binary string data type. For example, the following pairs of definitions are equivalent:

```
CHAR(10) CHARACTER SET binary
BINARY(10)

VARCHAR(10) CHARACTER SET binary
VARBINARY(10)

TEXT CHARACTER SET binary
BLOB
```

- `CAST(expr AS type)`

  The `CAST()` function takes an expression of any type and produces a result value of the specified type, similar to `CONVERT()`. For more information, see the description of `CONVERT()`.

  `CAST()` is standard SQL syntax.

- `CONVERT(expr,type)`, `CONVERT(expr USING transcoding_name)`

  The `CONVERT()` function takes an expression of any type and produces a result value of the specified type.

  Discussion of `CONVERT(expr, type)` syntax here also applies to `CAST(expr AS type)`, which is equivalent.

  `CONVERT(... USING ...)` is standard SQL syntax. The non-`USING` form of `CONVERT()` is ODBC syntax.

`CONVERT()` with `USING` converts data between different character sets. In MySQL, transcoding names are the same as the corresponding character set names. For example, this statement converts the string `'abc'` in the default character set to the corresponding string in the `utf8` character set:

```
SELECT CONVERT('abc' USING utf8);
```

`CONVERT()` without `USING` and `CAST()` take an expression and a *`type`* value specifying the result type. These *`type`* values are permitted:

- `BINARY[(N)]`

  Produces a string with the `BINARY` data type. See Section 11.4.2, “The BINARY and VARBINARY Types” for a description of how this affects comparisons. If the optional length *`N`* is given, `BINARY(N)` causes the cast to use no more than *`N`* bytes of the argument. Values shorter than *`N`* bytes are padded with `0x00` bytes to a length of *`N`*.

- `CHAR[(N)] [charset_info]`

  Produces a string with the `CHAR` data type. If the optional length *`N`* is given, `CHAR(N)` causes the cast to use no more than *`N`* characters of the argument. No padding occurs for values shorter than *`N`* characters.

  With no *`charset_info`* clause, `CHAR` produces a string with the default character set. To specify the character set explicitly, these *`charset_info`* values are permitted:

  - `CHARACTER SET charset_name`: Produces a string with the given character set.
  - `ASCII`: Shorthand for `CHARACTER SET latin1`.
  - `UNICODE`: Shorthand for `CHARACTER SET ucs2`.

  In all cases, the string has the default collation for the character set.

- `DATE`

  Produces a `DATE` value.

- `DATETIME`

  Produces a `DATETIME` value.

- `DECIMAL[(M[,D])]`

  Produces a `DECIMAL` value. If the optional *`M`* and *`D`* values are given, they specify the maximum number of digits (the precision) and the number of digits following the decimal point (the scale).

- `JSON`

  Produces a `JSON` value. For details on the rules for conversion of values between `JSON` and other types, see Comparison and Ordering of JSON Values.

- `NCHAR[(N)]`

  Like `CHAR`, but produces a string with the national character set. See Section 10.1.3.7, “The National Character Set”.

  Unlike `CHAR`, `NCHAR` does not permit trailing character set information to be specified.

- `SIGNED [INTEGER]`

Produces a signed integer value.

- `TIME`

  Produces a `TIME` value.

- `UNSIGNED [INTEGER]`

  Produces an unsigned integer value.

12.11 XML Functions

Table 12.15 XML Functions

Name	Description
`ExtractValue()`	Extracts a value from an XML string using XPath notation
`UpdateXML()`	Return replaced XML fragment

This section discusses XML and related functionality in MySQL.

Note

It is possible to obtain XML-formatted output from MySQL in the `mysql` and `mysqldump` clients by invoking them with the `--xml` option. See Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”, and Section 4.5.4, “`mysqldump` — A Database Backup Program”.

Two functions providing basic XPath 1.0 (XML Path Language, version 1.0) capabilities are available. Some basic information about XPath syntax and usage is provided later in this section; however, an in-depth discussion of these topics is beyond the scope of this Manual, and you should refer to the XML Path Language (XPath) 1.0 standard for definitive information. A useful resource for those new to XPath or who desire a refresher in the basics is the Zvon.org XPath Tutorial, which is available in several languages.

Note

These functions remain under development. We continue to improve these and other aspects of XML and XPath functionality in MySQL 8.0 and onwards. You may discuss these, ask questions about them, and obtain help from other users with them in the MySQL XML User Forum.

XPath expressions used with these functions support user variables and local stored program variables. User variables are weakly checked; variables local to stored programs are strongly checked (see also Bug #26518):

- **User variables (weak checking).** Variables using the syntax `$@variable_name` (that is, user variables) are not checked. No warnings or errors are issued by the server if a variable has the wrong type or has previously not been assigned a value. This also means the user is fully responsible for any typographical errors, since no warnings will be given if (for example) `$@myvariable` is used where `$@myvariable` was intended.

  Example:

```
mysql> SET @xml = '<a><b>X</b><b>Y</b></a>';
Query OK, 0 rows affected (0.00 sec)

mysql> SET @i =1, @j = 2;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @i, ExtractValue(@xml, '//b[$@i]');
+------+-------------------------------+
```

```
| @i   | ExtractValue(@xml, '//b[$@i]') |
+------+--------------------------------+
|    1 | X                              |
+------+--------------------------------+
1 row in set (0.00 sec)

mysql> SELECT @j, ExtractValue(@xml, '//b[$@j]');
+------+--------------------------------+
| @j   | ExtractValue(@xml, '//b[$@j]') |
+------+--------------------------------+
|    2 | Y                              |
+------+--------------------------------+
1 row in set (0.00 sec)

mysql> SELECT @k, ExtractValue(@xml, '//b[$@k]');
+------+--------------------------------+
| @k   | ExtractValue(@xml, '//b[$@k]') |
+------+--------------------------------+
| NULL |                                |
+------+--------------------------------+
1 row in set (0.00 sec)
```

- **Variables in stored programs (strong checking).** Variables using the syntax *`$variable_name`* can be declared and used with these functions when they are called inside stored programs. Such variables are local to the stored program in which they are defined, and are strongly checked for type and value.

  Example:

```
mysql> DELIMITER |

mysql> CREATE PROCEDURE myproc ()
    -> BEGIN
    ->   DECLARE i INT DEFAULT 1;
    ->   DECLARE xml VARCHAR(25) DEFAULT '<a>X</a><a>Y</a><a>Z</a>';
    ->
    ->   WHILE i < 4 DO
    ->     SELECT xml, i, ExtractValue(xml, '//a[$i]');
    ->     SET i = i+1;
    ->   END WHILE;
    -> END |
Query OK, 0 rows affected (0.01 sec)

mysql> DELIMITER ;

mysql> CALL myproc();
+-------------------------+---+------------------------------+
| xml                     | i | ExtractValue(xml, '//a[$i]') |
+-------------------------+---+------------------------------+
| <a>X</a><a>Y</a><a>Z</a> | 1 | X                            |
+-------------------------+---+------------------------------+
1 row in set (0.00 sec)

+-------------------------+---+------------------------------+
| xml                     | i | ExtractValue(xml, '//a[$i]') |
+-------------------------+---+------------------------------+
| <a>X</a><a>Y</a><a>Z</a> | 2 | Y                            |
+-------------------------+---+------------------------------+
1 row in set (0.01 sec)

+-------------------------+---+------------------------------+
| xml                     | i | ExtractValue(xml, '//a[$i]') |
+-------------------------+---+------------------------------+
| <a>X</a><a>Y</a><a>Z</a> | 3 | Z                            |
+-------------------------+---+------------------------------+
1 row in set (0.01 sec)
```

  Parameters. Variables used in XPath expressions inside stored routines that are passed in as parameters are also subject to strong checking.

Expressions containing user variables or variables local to stored programs must otherwise (except for notation) conform to the rules for XPath expressions containing variables as given in the XPath 1.0 specification.

Note

A user variable used to store an XPath expression is treated as an empty string. Because of this, it is not possible to store an XPath expression as a user variable. (Bug #32911)

- `ExtractValue(xml_frag, xpath_expr)`

  `ExtractValue()` takes two string arguments, a fragment of XML markup *`xml_frag`* and an XPath expression *`xpath_expr`* (also known as a *locator*); it returns the text (`CDATA`) of the first text node which is a child of the elements or elements matched by the XPath expression.

  Using this function is the equivalent of performing a match using the *`xpath_expr`* after appending `/text()`. In other words, `ExtractValue('<a><b>Sakila</b></a>', '/a/b')` and `ExtractValue('<a><b>Sakila</b></a>', '/a/b/text()')` produce the same result.

  If multiple matches are found, the content of the first child text node of each matching element is returned (in the order matched) as a single, space-delimited string.

  If no matching text node is found for the expression (including the implicit `/text()`)—for whatever reason, as long as *`xpath_expr`* is valid, and *`xml_frag`* consists of elements which are properly nested and closed—an empty string is returned. No distinction is made between a match on an empty element and no match at all. This is by design.

  If you need to determine whether no matching element was found in *`xml_frag`* or such an element was found but contained no child text nodes, you should test the result of an expression that uses the XPath `count()` function. For example, both of these statements return an empty string, as shown here:

```
mysql> SELECT ExtractValue('<a><b/></a>', '/a/b');
+-------------------------------------+
| ExtractValue('<a><b/></a>', '/a/b') |
+-------------------------------------+
|                                     |
+-------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT ExtractValue('<a><c/></a>', '/a/b');
+-------------------------------------+
| ExtractValue('<a><c/></a>', '/a/b') |
+-------------------------------------+
|                                     |
+-------------------------------------+
1 row in set (0.00 sec)
```

  However, you can determine whether there was actually a matching element using the following:

```
mysql> SELECT ExtractValue('<a><b/></a>', 'count(/a/b)');
+-------------------------------------+
| ExtractValue('<a><b/></a>', 'count(/a/b)') |
+-------------------------------------+
| 1                                   |
+-------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT ExtractValue('<a><c/></a>', 'count(/a/b)');
+-------------------------------------+
| ExtractValue('<a><c/></a>', 'count(/a/b)') |
+-------------------------------------+
| 0                                   |
```

```
+-----------------------------------+
1 row in set (0.01 sec)
```

Important

`ExtractValue()` returns only `CDATA`, and does not return any tags that might be contained within a matching tag, nor any of their content (see the result returned as `val1` in the following example).

```
mysql> SELECT
    ->   ExtractValue('<a>ccc<b>ddd</b></a>', '/a') AS val1,
    ->   ExtractValue('<a>ccc<b>ddd</b></a>', '/a/b') AS val2,
    ->   ExtractValue('<a>ccc<b>ddd</b></a>', '//b') AS val3,
    ->   ExtractValue('<a>ccc<b>ddd</b></a>', '/b') AS val4,
    ->   ExtractValue('<a>ccc<b>ddd</b><b>eee</b></a>', '//b') AS val5;

+------+------+------+------+---------+
| val1 | val2 | val3 | val4 | val5    |
+------+------+------+------+---------+
| ccc  | ddd  | ddd  |      | ddd eee |
+------+------+------+------+---------+
```

This function uses the current SQL collation for making comparisons with `contains()`, performing the same collation aggregation as other string functions (such as `CONCAT()`), in taking into account the collation coercibility of their arguments; see Section 10.1.8.4, “Collation Coercibility in Expressions”, for an explanation of the rules governing this behavior.

(Previously, binary—that is, case-sensitive—comparison was always used.)

`NULL` is returned if *`xml_frag`* contains elements which are not properly nested or closed, and a warning is generated, as shown in this example:

```
mysql> SELECT ExtractValue('<a>c</a><b', '//a');
+-----------------------------------+
| ExtractValue('<a>c</a><b', '//a') |
+-----------------------------------+
| NULL                              |
+-----------------------------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1525
Message: Incorrect XML value: 'parse error at line 1 pos 11:
         END-OF-INPUT unexpected ('>' wanted)'
1 row in set (0.00 sec)

mysql> SELECT ExtractValue('<a>c</a><b/>', '//a');
+-------------------------------------+
| ExtractValue('<a>c</a><b/>', '//a') |
+-------------------------------------+
| c                                   |
+-------------------------------------+
1 row in set (0.00 sec)
```

- `UpdateXML(`*`xml_target`*`, `*`xpath_expr`*`, `*`new_xml`*`)`

  This function replaces a single portion of a given fragment of XML markup *`xml_target`* with a new XML fragment *`new_xml`*, and then returns the changed XML. The portion of *`xml_target`* that is replaced matches an XPath expression *`xpath_expr`* supplied by the user.

  If no expression matching *`xpath_expr`* is found, or if multiple matches are found, the function returns the original *`xml_target`* XML fragment. All three arguments should be strings.

```
mysql> SELECT
    ->   UpdateXML('<a><b>ccc</b><d></d></a>', '/a', '<e>fff</e>') AS val1,
    ->   UpdateXML('<a><b>ccc</b><d></d></a>', '/b', '<e>fff</e>') AS val2,
    ->   UpdateXML('<a><b>ccc</b><d></d></a>', '//b', '<e>fff</e>') AS val3,
    ->   UpdateXML('<a><b>ccc</b><d></d></a>', '/a/d', '<e>fff</e>') AS val4,
    ->   UpdateXML('<a><d></d><b>ccc</b><d></d></a>', '/a/d', '<e>fff</e>') AS val5
    -> \G

*************************** 1. row ***************************
val1: <e>fff</e>
val2: <a><b>ccc</b><d></d></a>
val3: <a><e>fff</e><d></d></a>
val4: <a><b>ccc</b><e>fff</e></a>
val5: <a><d></d><b>ccc</b><d></d></a>
```

Note

A discussion in depth of XPath syntax and usage are beyond the scope of this Manual. Please see the XML Path Language (XPath) 1.0 specification for definitive information. A useful resource for those new to XPath or who are wishing a refresher in the basics is the Zvon.org XPath Tutorial, which is available in several languages.

Descriptions and examples of some basic XPath expressions follow:

- */tag*

  Matches *<tag/>* if and only if *<tag/>* is the root element.

  Example: `/a` has a match in `<a><b/></a>` because it matches the outermost (root) tag. It does not match the inner *a* element in `<b><a/></b>` because in this instance it is the child of another element.

- */tag1/tag2*

  Matches *<tag2/>* if and only if it is a child of *<tag1/>*, and *<tag1/>* is the root element.

  Example: `/a/b` matches the *b* element in the XML fragment `<a><b/></a>` because it is a child of the root element *a*. It does not have a match in `<b><a/></b>` because in this case, *b* is the root element (and hence the child of no other element). Nor does the XPath expression have a match in `<a><c><b/></c></a>`; here, *b* is a descendant of *a*, but not actually a child of *a*.

  This construct is extendable to three or more elements. For example, the XPath expression `/a/b/c` matches the *c* element in the fragment `<a><b><c/></b></a>`.

- *//tag*

  Matches any instance of *<tag>*.

  Example: `//a` matches the *a* element in any of the following: `<a><b><c/></b></a>`; `<c><a><b/></a></b>`; `<c><b><a/></b></c>`.

  `//` can be combined with `/`. For example, `//a/b` matches the *b* element in either of the fragments `<a><b/></a>` or `<a><b><c/></b></a>`

Note

//tag is the equivalent of `/descendant-or-self::*/tag`. A common error is to confuse this with `/descendant-or-self::tag`, although the latter expression can actually lead to very different results, as can be seen here:

```
mysql> SET @xml = '<a><b><c>w</c><b>x</b><d>y</d>z</b></a>';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @xml;
+---------------------------------------+
| @xml                                  |
+---------------------------------------+
| <a><b><c>w</c><b>x</b><d>y</d>z</b></a> |
+---------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT ExtractValue(@xml, '//b[1]');
+-------------------------------+
| ExtractValue(@xml, '//b[1]') |
+-------------------------------+
| x z                           |
+-------------------------------+
1 row in set (0.00 sec)

mysql> SELECT ExtractValue(@xml, '//b[2]');
+-------------------------------+
| ExtractValue(@xml, '//b[2]') |
+-------------------------------+
|                               |
+-------------------------------+
1 row in set (0.01 sec)

mysql> SELECT ExtractValue(@xml, '/descendant-or-self::*/b[1]');
+-------------------------------------------------+
| ExtractValue(@xml, '/descendant-or-self::*/b[1]') |
+-------------------------------------------------+
| x z                                             |
+-------------------------------------------------+
1 row in set (0.06 sec)

mysql> SELECT ExtractValue(@xml, '/descendant-or-self::*/b[2]');
+-------------------------------------------------+
| ExtractValue(@xml, '/descendant-or-self::*/b[2]') |
+-------------------------------------------------+
|                                                 |
+-------------------------------------------------+
1 row in set (0.00 sec)


mysql> SELECT ExtractValue(@xml, '/descendant-or-self::b[1]');
+-----------------------------------------------+
| ExtractValue(@xml, '/descendant-or-self::b[1]') |
+-----------------------------------------------+
| z                                             |
+-----------------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT ExtractValue(@xml, '/descendant-or-self::b[2]');
+-----------------------------------------------+
| ExtractValue(@xml, '/descendant-or-self::b[2]') |
+-----------------------------------------------+
| x                                             |
+-----------------------------------------------+
1 row in set (0.00 sec)
```

- The `*` operator acts as a “wildcard” that matches any element. For example, the expression `/*/b` matches the *b* element in either of the XML fragments `<a><b/></a>` or `<c><b/></c>`. However, the expression does not produce a match in the fragment `<b><a/></b>` because *b* must be a child of some other element. The wildcard may be used in any position: The expression `/*/b/*` will match any child of a *b* element that is itself not the root element.

- You can match any of several locators using the `|` (`UNION`) operator. For example, the expression `//b|//c` matches all *b* and *c* elements in the XML target.

- It is also possible to match an element based on the value of one or more of its attributes. This done using the syntax `tag[@attribute="value"]`. For example, the expression `//b[@id="idB"]` matches the second *b* element in the fragment `<a><b id="idA"/><c/><b id="idB"/></a>`. To match against *any* element having `attribute="value"`, use the XPath expression `//*[attribute="value"]`.

  To filter multiple attribute values, simply use multiple attribute-comparison clauses in succession. For example, the expression `//b[@c="x"][@d="y"]` matches the element `<b c="x" d="y"/>` occurring anywhere in a given XML fragment.

  To find elements for which the same attribute matches any of several values, you can use multiple locators joined by the `|` operator. For example, to match all *b* elements whose *c* attributes have either of the values 23 or 17, use the expression `//b[@c="23"]|//b[@c="17"]`. You can also use the logical `or` operator for this purpose: `//b[@c="23" or @c="17"]`.

  Note

  The difference between `or` and `|` is that `or` joins conditions, while `|` joins result sets.

XPath Limitations. The XPath syntax supported by these functions is currently subject to the following limitations:

- Nodeset-to-nodeset comparison (such as `'/a/b[@c=@d]'`) is not supported.
- All of the standard XPath comparison operators are supported. (Bug #22823)
- Relative locator expressions are resolved in the context of the root node. For example, consider the following query and result:

```
mysql> SELECT ExtractValue(
    ->   '<a><b c="1">X</b><b c="2">Y</b></a>',
    ->    'a/b'
    -> ) AS result;
+--------+
| result |
+--------+
| X Y    |
+--------+
1 row in set (0.03 sec)
```

  In this case, the locator `a/b` resolves to `/a/b`.

  Relative locators are also supported within predicates. In the following example, `d[../@c="1"]` is resolved as `/a/b[@c="1"]/d`:

```
mysql> SELECT ExtractValue(
    ->      '<a>
    ->        <b c="1"><d>X</d></b>
    ->        <b c="2"><d>X</d></b>
    ->      </a>',
    ->      'a/b/d[../@c="1"]')
    -> AS result;
+--------+
| result |
+--------+
| X      |
+--------+
1 row in set (0.00 sec)
```

- Locators prefixed with expressions that evaluate as scalar values—including variable references, literals, numbers, and scalar function calls—are not permitted, and their use results in an error.
- The `::` operator is not supported in combination with node types such as the following:

- *axis*::comment()
- *axis*::text()
- *axis*::processing-instructions()
- *axis*::node()

However, name tests (such as *axis*::*name* and *axis*::*) are supported, as shown in these examples:

```
mysql> SELECT ExtractValue('<a><b>x</b><c>y</c></a>','/a/child::b');
+-----------------------------------------------------+
| ExtractValue('<a><b>x</b><c>y</c></a>','/a/child::b') |
+-----------------------------------------------------+
| x                                                   |
+-----------------------------------------------------+
1 row in set (0.02 sec)

mysql> SELECT ExtractValue('<a><b>x</b><c>y</c></a>','/a/child::*');
+-----------------------------------------------------+
| ExtractValue('<a><b>x</b><c>y</c></a>','/a/child::*') |
+-----------------------------------------------------+
| x y                                                 |
+-----------------------------------------------------+
1 row in set (0.01 sec)
```

- “Up-and-down” navigation is not supported in cases where the path would lead “above” the root element. That is, you cannot use expressions which match on descendants of ancestors of a given element, where one or more of the ancestors of the current element is also an ancestor of the root element (see Bug #16321).
- The following XPath functions are not supported, or have known issues as indicated:
  - id()
  - lang()
  - local-name()
  - name()
  - namespace-uri()
  - normalize-space()
  - starts-with()
  - string()
  - substring-after()
  - substring-before()
  - translate()
- The following axes are not supported:
  - following-sibling
  - following
  - preceding-sibling

- `preceding`

XPath expressions passed as arguments to `ExtractValue()` and `UpdateXML()` may contain the colon character (`:`) in element selectors, which enables their use with markup employing XML namespaces notation. For example:

```
mysql> SET @xml = '<a>111<b:c>222<d>333</d><e:f>444</e:f></b:c></a>';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT ExtractValue(@xml, '//e:f');
+-----------------------------+
| ExtractValue(@xml, '//e:f') |
+-----------------------------+
| 444                         |
+-----------------------------+
1 row in set (0.00 sec)

mysql> SELECT UpdateXML(@xml, '//b:c', '<g:h>555</g:h>');
+-------------------------------------------+
| UpdateXML(@xml, '//b:c', '<g:h>555</g:h>') |
+-------------------------------------------+
| <a>111<g:h>555</g:h></a>                  |
+-------------------------------------------+
1 row in set (0.00 sec)
```

This is similar in some respects to what is permitted by Apache Xalan and some other parsers, and is much simpler than requiring namespace declarations or the use of the `namespace-uri()` and `local-name()` functions.

Error handling. For both `ExtractValue()` and `UpdateXML()`, the XPath locator used must be valid and the XML to be searched must consist of elements which are properly nested and closed. If the locator is invalid, an error is generated:

```
mysql> SELECT ExtractValue('<a>c</a><b/>', '/&a');
ERROR 1105 (HY000): XPATH syntax error: '&a'
```

If *`xml_frag`* does not consist of elements which are properly nested and closed, `NULL` is returned and a warning is generated, as shown in this example:

```
mysql> SELECT ExtractValue('<a>c</a><b', '//a');
+----------------------------------+
| ExtractValue('<a>c</a><b', '//a') |
+----------------------------------+
| NULL                             |
+----------------------------------+
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1525
Message: Incorrect XML value: 'parse error at line 1 pos 11:
         END-OF-INPUT unexpected ('>' wanted)'
1 row in set (0.00 sec)

mysql> SELECT ExtractValue('<a>c</a><b/>', '//a');
+-----------------------------------+
| ExtractValue('<a>c</a><b/>', '//a') |
+-----------------------------------+
| c                                 |
+-----------------------------------+
1 row in set (0.00 sec)
```

Important

The replacement XML used as the third argument to `UpdateXML()` is *not* checked to determine whether it consists solely of elements which are properly nested and closed.

XPath Injection. *code injection* occurs when malicious code is introduced into the system to gain unauthorized access to privileges and data. It is based on exploiting assumptions made by developers about the type and content of data input from users. XPath is no exception in this regard.

A common scenario in which this can happen is the case of application which handles authorization by matching the combination of a login name and password with those found in an XML file, using an XPath expression like this one:

```
//user[login/text()='neapolitan' and password/text()='1c3cr34m']/attribute::id
```

This is the XPath equivalent of an SQL statement like this one:

```
SELECT id FROM users WHERE login='neapolitan' AND password='1c3cr34m';
```

A PHP application employing XPath might handle the login process like this:

```
<?php

  $file     =   "users.xml";

  $login    =   $POST["login"];
  $password =   $POST["password"];

  $xpath = "//user[login/text()=$login and password/text()=$password]/attribute::id";

  if( file_exists($file) )
  {
    $xml = simplexml_load_file($file);

    if($result = $xml->xpath($xpath))
      echo "You are now logged in as user $result[0].";
    else
      echo "Invalid login name or password.";
  }
  else
    exit("Failed to open $file.");

?>
```

No checks are performed on the input. This means that a malevolent user can “short-circuit” the test by entering `' or 1=1` for both the login name and password, resulting in `$xpath` being evaluated as shown here:

```
//user[login/text()='' or 1=1 and password/text()='' or 1=1]/attribute::id
```

Since the expression inside the square brackets always evaluates as `true`, it is effectively the same as this one, which matches the `id` attribute of every `user` element in the XML document:

```
//user/attribute::id
```

One way in which this particular attack can be circumvented is simply by quoting the variable names to be interpolated in the definition of `$xpath`, forcing the values passed from a Web form to be converted to strings:

```
$xpath = "//user[login/text()='$login' and password/text()='$password']/attribute::id";
```

This is the same strategy that is often recommended for preventing SQL injection attacks. In general, the practices you should follow for preventing XPath injection attacks are the same as for preventing SQL injection:

- Never accepted untested data from users in your application.
- Check all user-submitted data for type; reject or convert data that is of the wrong type
- Test numeric data for out of range values; truncate, round, or reject values that are out of range. Test strings for illegal characters and either strip them out or reject input containing them.
- Do not output explicit error messages that might provide an unauthorized user with clues that could be used to compromise the system; log these to a file or database table instead.

Just as SQL injection attacks can be used to obtain information about database schemas, so can XPath injection be used to traverse XML files to uncover their structure, as discussed in Amit Klein's paper Blind XPath Injection (PDF file, 46KB).

It is also important to check the output being sent back to the client. Consider what can happen when we use the MySQL `ExtractValue()` function:

```
mysql> SELECT ExtractValue(
    ->     LOAD_FILE('users.xml'),
    ->     '//user[login/text()="" or 1=1 and password/text()="" or 1=1]/attribute::id'
    -> ) AS id;
+-------------------------------+
| id                            |
+-------------------------------+
| 00327 13579 02403 42354 28570 |
+-------------------------------+
1 row in set (0.01 sec)
```

Because `ExtractValue()` returns multiple matches as a single space-delimited string, this injection attack provides every valid ID contained within `users.xml` to the user as a single row of output. As an extra safeguard, you should also test output before returning it to the user. Here is a simple example:

```
mysql> SELECT @id = ExtractValue(
    ->     LOAD_FILE('users.xml'),
    ->     '//user[login/text()="" or 1=1 and password/text()="" or 1=1]/attribute::id'
    -> );
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT IF(
    ->     INSTR(@id, ' ') = 0,
    ->     @id,
    ->     'Unable to retrieve user ID')
    -> AS singleID;
+---------------------------+
| singleID                  |
+---------------------------+
| Unable to retrieve user ID |
+---------------------------+
1 row in set (0.00 sec)
```

In general, the guidelines for returning data to users securely are the same as for accepting user input. These can be summed up as:

- Always test outgoing data for type and permissible values.
- Never permit unauthorized users to view error messages that might provide information about the application that could be used to exploit it.

12.12 Bit Functions and Operators

Table 12.16 Bit Functions and Operators

Name	Description
`BIT_COUNT()`	Return the number of bits that are set
`&`	Bitwise AND
`~`	Bitwise inversion
`|`	Bitwise OR
`^`	Bitwise XOR
`<<`	Left shift
`>>`	Right shift

Bit functions and operators comprise `BIT_COUNT()`, `BIT_AND()`, `BIT_OR()`, `BIT_XOR()`, `&`, `|`, `^`, `~`, `<<`, and `>>`. (The `BIT_AND()`, `BIT_OR()`, and `BIT_XOR()` aggregate functions are described in Section 12.18.1, “Aggregate (GROUP BY) Function Descriptions”.) Prior to MySQL 8.0, bit functions and operators required `BIGINT` (64-bit integer) arguments and returned `BIGINT` values, so they had a maximum range of 64 bits. Non-`BIGINT` arguments were converted to `BIGINT` prior to performing the operation and truncation could occur.

In MySQL 8.0, bit functions and operators permit binary string type arguments (`BINARY`, `VARBINARY`, and the `BLOB` types) and return a value of like type, which enables them to take arguments and produce return values larger than 64 bits. Nonbinary string arguments are converted to `BIGINT` and processed as such, as before.

An implication of this change in behavior is that bit operations on binary string arguments might produce a different result in MySQL 8.0 than in 5.7. For information about how to prepare in MySQL 5.7 for potential incompatibilities between MySQL 5.7 and 8.0, see Bit Functions and Operators, in MySQL 5.7 Reference Manual.

- Bit Operations Prior to MySQL 8.0
- Bit Operations in MySQL 8.0
- Binary String Bit-Operation Examples
- Bitwise AND, OR, and XOR Operations
- Bitwise Complement and Shift Operations
- BIT_COUNT() Operations
- BIT_AND(), BIT_OR(), and BIT_XOR() Operations
- Special Handling of Hexadecimal Literals, Bit Literals, and NULL Literals
- Bit-Operation Incompatibilities with MySQL 5.7

Bit Operations Prior to MySQL 8.0

Bit operations prior to MySQL 8.0 handle only unsigned 64-bit integer argument and result values (that is, unsigned `BIGINT` values). Conversion of arguments of other types to `BIGINT` occurs as necessary. Examples:

- This statement operates on numeric literals, treated as unsigned 64-bit integers:

```
mysql> SELECT 127 | 128, 128 << 2, BIT_COUNT(15);
+-----------+----------+---------------+
| 127 | 128 | 128 << 2 | BIT_COUNT(15) |
+-----------+----------+---------------+
|       255 |      512 |             4 |
+-----------+----------+---------------+
```

- This statement performs to-number conversions on the string arguments (`'127'` to `127`, and so forth) before performing the same operations as the first statement and producing the same results:

```
mysql> SELECT '127' | '128', '128' << 2, BIT_COUNT('15');
+--------------+------------+----------------+
| '127' | '128' | '128' << 2 | BIT_COUNT('15') |
+--------------+------------+----------------+
|          255 |        512 |              4 |
+--------------+------------+----------------+
```

- This statement uses hexadecimal literals for the bit-operation arguments. MySQL by default treats hexadecimal literals as binary strings, but in numeric context evaluates them as numbers (see Section 9.1.4, “Hexadecimal Literals”). Prior to MySQL 8.0, numeric context includes bit operations. Examples:

```
mysql> SELECT X'7F' | X'80', X'80' << 2, BIT_COUNT(X'0F');
+--------------+------------+-----------------+
| X'7F' | X'80' | X'80' << 2 | BIT_COUNT(X'0F') |
+--------------+------------+-----------------+
|          255 |        512 |               4 |
+--------------+------------+-----------------+
```

  Handling of bit-value literals in bit operations is similar to hexadecimal literals (that is, as numbers).

Bit Operations in MySQL 8.0

MySQL 8.0 extends bit operations to handle binary string arguments directly (without conversion) and produce binary string results. (Arguments that are not integers or binary strings are still converted to integers, as before.) This extension enhances bit operations in the following ways:

- Bit operations become possible on values longer than 64 bits.
- It is easier to perform bit operations on values that are more naturally represented as binary strings than as integers.

For example, consider UUID values and IPv6 addresses, which have human-readable text formats like this:

```
UUID: 6ccd780c-baba-1026-9564-5b8c656024db
IPv6: fe80::219:d1ff:fe91:1a72
```

It is cumbersome to operate on text strings in those formats. An alternative is convert them to fixed-length binary strings without delimiters. `UUID_TO_BIN()` and `INET6_ATON()` each produce a value of data type `BINARY(16)`, a binary string 16 bytes (128 bits) long. The following statements illustrate this (`HEX()` is used to produce displayable values):

```
mysql> SELECT HEX(UUID_TO_BIN('6ccd780c-baba-1026-9564-5b8c656024db'));
+--------------------------------------------------------+
| HEX(UUID_TO_BIN('6ccd780c-baba-1026-9564-5b8c656024db')) |
+--------------------------------------------------------+
| 6CCD780CBABA102695645B8C656024DB                       |
+--------------------------------------------------------+
mysql> SELECT HEX(INET6_ATON('fe80::219:d1ff:fe91:1a72'));
+-------------------------------------------+
| HEX(INET6_ATON('fe80::219:d1ff:fe91:1a72')) |
+-------------------------------------------+
| FE800000000000000219D1FFFE911A72          |
+-------------------------------------------+
```

Those binary values are easily manipulable with bit operations to perform actions such as extracting the timestamp from UUID values, or extracting the network and host parts of IPv6 addresses. (For examples, see later in this discussion.)

Arguments that count as binary strings include column values, routine parameters, local variables, and user-defined variables that have a binary string type: `BINARY`, `VARBINARY`, or one of the `BLOB` types.

What about hexadecimal literals and bit literals? Recall that those are binary strings by default in MySQL, but numbers in numeric context. How are they handled for bit operations in MySQL 8.0? Does MySQL continue to evaluate them in numeric context, as is done prior to MySQL 8.0? Or do bit operations evaluate them as binary strings, now that binary strings can be handled “natively” without conversion?

Answer: It has been common to specify arguments to bit operations using hexadecimal literals or bit literals with the intent that they represent numbers, so MySQL continues to evaluate bit operations in numeric context when all bit arguments are hexadecimal or bit literals, for backward compatility. If you require evaluation as binary strings instead, that is easily accomplished: Use the `_binary` introducer for at least one literal.

- These bit operations evaluate the hexadecimal literals and bit literals as integers:

```
mysql> SELECT X'40' | X'01', b'11110001' & b'01001111';
+--------------+--------------------------+
| X'40' | X'01' | b'11110001' & b'01001111' |
+--------------+--------------------------+
|           65 |                       65 |
+--------------+--------------------------+
```

- These bit operations evaluate the hexadecimal literals and bit literals as binary strings, due to the `_binary` introducer:

```
mysql> SELECT _binary X'40' | X'01', b'11110001' & _binary b'01001111';
+----------------------+----------------------------------+
| _binary X'40' | X'01' | b'11110001' & _binary b'01001111' |
+----------------------+----------------------------------+
| A                    | A                                |
+----------------------+----------------------------------+
```

Although the bit operations in both statements produce a result with a numeric value of 65, the second statement operates in binary-string context, for which 65 is ASCII `A`.

In numeric evaluation context, permitted values of hexadecimal literal and bit literal arguments have a maximum of 64 bits, as do results. By contrast, in binary-string evaluation context, permitted arguments (and results) can exceed 64 bits:

```
mysql> SELECT _binary X'4040404040404040' | X'0102030405060708';
+--------------------------------------------------+
| _binary X'4040404040404040' | X'0102030405060708' |
+--------------------------------------------------+
| ABCDEFGH                                         |
+--------------------------------------------------+
```

There are several ways to refer to a hexadecimal literal or bit literal in a bit operation to cause binary-string evaluation:

```
_binary literal
BINARY literal
CAST(literal AS BINARY)
```

Another way to produce binary-string evaluation of hexadecimal literals or bit literals is to assign them to user-defined variables, which results in variables that have a binary string type:

```
mysql> SET @v1 = X'40', @v2 = X'01', @v3 = b'11110001', @v4 = b'01001111';
mysql> SELECT @v1 | @v2, @v3 & @v4;
+-----------+-----------+
| @v1 | @v2 | @v3 & @v4 |
+-----------+-----------+
```

```
| A         | A         |
+-----------+-----------+
```

In binary-string context, bitwise operation arguments must have the same length or an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs:

```
mysql> SELECT _binary X'40' | X'0001';
ERROR 3513 (HY000): Binary operands of bitwise
operators must be of equal length
```

To satisfy the equal-length requirement, pad the shorter value with leading zero digits or, if the longer value begins with leading zero digits and a shorter result value is acceptable, strip them:

```
mysql> SELECT _binary X'0040' | X'0001';
+--------------------------+
| _binary X'0040' | X'0001' |
+--------------------------+
|  A                       |
+--------------------------+
mysql> SELECT _binary X'40' | X'01';
+----------------------+
| _binary X'40' | X'01' |
+----------------------+
| A                    |
+----------------------+
```

Padding or stripping can also be accomplished using functions such as `LPAD()`, `RPAD()`, `SUBSTR()`, or `CAST()`. In such cases, the expression arguments are no longer all literals and `_binary` becomes unnecessary. Examples:

```
mysql> SELECT LPAD(X'40', 2, X'00') | X'0001';
+---------------------------------+
| LPAD(X'40', 2, X'00') | X'0001' |
+---------------------------------+
|  A                              |
+---------------------------------+
mysql> SELECT X'40' | SUBSTR(X'0001', 2, 1);
+-----------------------------+
| X'40' | SUBSTR(X'0001', 2, 1) |
+-----------------------------+
| A                           |
+-----------------------------+
```

Binary String Bit-Operation Examples

The following example illustrates use of bit operations to extract parts of a UUID value, in this case, the timestamp and IEEE 802 node number. This technique requires bitmasks for each extracted part.

Convert the text UUID to the corresponding 16-byte binary value so that it can be manipulated using bit operations in binary-string context:

```
mysql> SET @uuid = UUID_TO_BIN('6ccd780c-baba-1026-9564-5b8c656024db');
mysql> SELECT HEX(@uuid);
+----------------------------------+
| HEX(@uuid)                       |
+----------------------------------+
| 6CCD780CBABA102695645B8C656024DB |
+----------------------------------+
```

Construct bitmasks for the timestamp and node number parts of the value. The timestamp comprises the first three parts (64 bits, bits 0 to 63) and the node number is the last part (48 bits, bits 80 to 127):

```
mysql> SET @ts_mask = CAST(X'FFFFFFFFFFFFFFFF' AS BINARY(16));
mysql> SET @node_mask = CAST(X'FFFFFFFFFFFF' AS BINARY(16)) >> 80;
mysql> SELECT HEX(@ts_mask);
```

```
+----------------------------------+
| HEX(@ts_mask)                    |
+----------------------------------+
| FFFFFFFFFFFFFFF0000000000000000 |
+----------------------------------+
mysql> SELECT HEX(@node_mask);
+----------------------------------+
| HEX(@node_mask)                  |
+----------------------------------+
| 00000000000000000000FFFFFFFFFFFF |
+----------------------------------+
```

The `CAST(... AS BINARY(16))` function is used here because the masks must be the same length as the UUID value against which they are applied. The same result can be produced using other functions to pad the masks to the required length:

```
SET @ts_mask= RPAD(X'FFFFFFFFFFFFFFF' , 16, X'00');
SET @node_mask = LPAD(X'FFFFFFFFFFFF', 16, X'00') ;
```

Use the masks to extract the timestamp and node number parts:

```
mysql> SELECT HEX(@uuid & @ts_mask) AS 'timestamp part';
+----------------------------------+
| timestamp part                   |
+----------------------------------+
| 6CCD780CBABA10260000000000000000 |
+----------------------------------+
mysql> SELECT HEX(@uuid & @node_mask) AS 'node part';
+----------------------------------+
| node part                        |
+----------------------------------+
| 00000000000000000005B8C656024DB |
+----------------------------------+
```

The preceding example uses these bit operations: right shift (`>>`) and bitwise AND (`&`).

Note

`UUID_TO_BIN()` takes a flag that causes some bit rearrangement in the resulting binary UUID value. If you use that flag, modify the extraction masks accordingly.

The next example uses bit operations to extract the network and host parts of an IPv6 address. Suppose that the network part has a length of 80 bits. Then the host part has a length of 128 − 80 = 48 bits. To extract the network and host parts of the address, convert it to a binary string, then use bit operations in binary-string context.

Convert the text IPv6 address to the corresponding binary string:

```
mysql> SET @ip = INET6_ATON('fe80::219:d1ff:fe91:1a72');
```

Define the network length in bits:

```
mysql> SET @net_len = 80;
```

Construct network and host masks by shifting the all-ones address left or right. To do this, begin with the address `::`, which is shorthand for all zeros, as you can see by converting it to a binary string like this:

```
mysql> SELECT HEX(INET6_ATON('::')) AS 'all zeros';
+----------------------------------+
| all zeros                        |
+----------------------------------+
| 00000000000000000000000000000000 |
```

```
+----------------------------------+
```

To produce the complementary value (all ones), use the ~ operator to invert the bits:

```
mysql> SELECT HEX(~INET6_ATON('::')) AS 'all ones';
+----------------------------------+
| all ones                         |
+----------------------------------+
| FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF |
+----------------------------------+
```

Shift the all-ones value left or right to produce the network and host masks:

```
mysql> SET @net_mask = ~INET6_ATON('::') << (128 - @net_len);
mysql> SET @host_mask = ~INET6_ATON('::') >> @net_len;
```

Display the masks to verify that they cover the correct parts of the address:

```
mysql> SELECT INET6_NTOA(@net_mask) AS 'network mask';
+-----------------------------+
| network mask                |
+-----------------------------+
| ffff:ffff:ffff:ffff:ffff:: |
+-----------------------------+
mysql> SELECT INET6_NTOA(@host_mask) AS 'host mask';
+------------------------+
| host mask              |
+------------------------+
| ::ffff:255.255.255.255 |
+------------------------+
```

Extract and display the network and host parts of the address:

```
mysql> SET @net_part = @ip & @net_mask;
mysql> SET @host_part = @ip & @host_mask;
mysql> SELECT INET6_NTOA(@net_part) AS 'network part';
+-----------------+
| network part    |
+-----------------+
| fe80::219:0:0:0 |
+-----------------+
mysql> SELECT INET6_NTOA(@host_part) AS 'host part';
+------------------+
| host part        |
+------------------+
| ::d1ff:fe91:1a72 |
+------------------+
```

The preceding example uses these bit operations: Complement (~), left shift (<<), and bitwise AND (&).

The remaining discussion provides details on argument handling for each group of bit operations, more information about literal-value handling in bit operations, and potential incompatibilities between MySQL 8.0 and older MySQL versions.

Bitwise AND, OR, and XOR Operations

For &, |, and ^ bit operations, the result type depends on whether the arguments are evaluated as binary strings or numbers:

- Binary-string evaluation occurs when the arguments have a binary string type, and at least one of them is not a hexadecimal literal, bit literal, or NULL literal. Numeric evaluation occurs otherwise, with argument conversion to unsigned 64-bit integers as necessary.
- Binary-string evaluation produces a binary string of the same length as the arguments. If the arguments have unequal lengths, an ER_INVALID_BITWISE_OPERANDS_SIZE error occurs. Numeric evaluation produces an unsigned 64-bit integer.

Examples of numeric evaluation:

```
mysql> SELECT 64 | 1, X'40' | X'01';
+--------+---------------+
| 64 | 1 | X'40' | X'01' |
+--------+---------------+
|     65 |            65 |
+--------+---------------+
```

Examples of binary-string evaluation:

```
mysql> SELECT _binary X'40' | X'01';
+-----------------------+
| _binary X'40' | X'01' |
+-----------------------+
| A                     |
+-----------------------+
mysql> SET @var1 = X'40', @var2 = X'01';
mysql> SELECT @var1 | @var2;
+---------------+
| @var1 | @var2 |
+---------------+
| A             |
+---------------+
```

Bitwise Complement and Shift Operations

For ~, <<, and >> bit operations, the result type depends on whether the bit argument is evaluated as a binary string or number:

- Binary-string evaluation occurs when the bit argument has a binary string type, and is not a hexadecimal literal, bit literal, or NULL literal. Numeric evaluation occurs otherwise, with argument conversion to an unsigned 64-bit integer as necessary.
- Binary-string evaluation produces a binary string of the same length as the bit argument. Numeric evaluation produces an unsigned 64-bit integer.

For shift operations, bits shifted off the end of the value are lost without warning, regardless of the argument type. In particular, if the shift count is greater or equal to the number of bits in the bit argument, all bits in the result are 0.

Examples of numeric evaluation:

```
mysql> SELECT ~0, 64 << 2, X'40' << 2;
+----------------------+---------+------------+
| ~0                   | 64 << 2 | X'40' << 2 |
+----------------------+---------+------------+
| 18446744073709551615 |     256 |        256 |
+----------------------+---------+------------+
```

Examples of binary-string evaluation:

```
mysql> SELECT HEX(_binary X'1111000022220000' >> 16);
+----------------------------------------+
| HEX(_binary X'1111000022220000' >> 16) |
+----------------------------------------+
| 0000111100002222                       |
+----------------------------------------+
mysql> SELECT HEX(_binary X'1111000022220000' << 16);
+----------------------------------------+
| HEX(_binary X'1111000022220000' << 16) |
+----------------------------------------+
| 0000222200000000                       |
+----------------------------------------+
mysql> SET @var1 = X'F0F0F0F0';
```

```
mysql> SELECT HEX(~@var1);
+-------------+
| HEX(~@var1) |
+-------------+
| 0F0F0F0F    |
+-------------+
```

BIT_COUNT() Operations

The `BIT_COUNT()` function always returns an unsigned 64-bit integer, or `NULL` if the argument is `NULL`.

```
mysql> SELECT BIT_COUNT(127);
+----------------+
| BIT_COUNT(127) |
+----------------+
|              7 |
+----------------+
mysql> SELECT BIT_COUNT(b'010101'), BIT_COUNT(_binary b'010101');
+----------------------+-----------------------------+
| BIT_COUNT(b'010101') | BIT_COUNT(_binary b'010101') |
+----------------------+-----------------------------+
|                    3 |                           3 |
+----------------------+-----------------------------+
```

BIT_AND(), BIT_OR(), and BIT_XOR() Operations

For the `BIT_AND()`, `BIT_OR()`, and `BIT_XOR()` bit functions, the result type depends on whether the function argument values are evaluated as binary strings or numbers:

- Binary-string evaluation occurs when the argument values have a binary string type, and the argument is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument value conversion to unsigned 64-bit integers as necessary.
- Binary-string evaluation produces a binary string of the same length as the argument values. If argument values have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. If the argument size exceeds 511 bytes, an `ER_INVALID_BITWISE_AGGREGATE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

`NULL` values do not affect the result unless all values are `NULL`. In that case, the result is a neutral value having the same length as the length of the argument values (all bits 1 for `BIT_AND()`, all bits 0 for `BIT_OR()`, and `BIT_XOR()`).

Example:

```
mysql> CREATE TABLE t (group_id INT, a VARBINARY(6));
mysql> INSERT INTO t VALUES (1, NULL);
mysql> INSERT INTO t VALUES (1, NULL);
mysql> INSERT INTO t VALUES (2, NULL);
mysql> INSERT INTO t VALUES (2, X'1234');
mysql> INSERT INTO t VALUES (2, X'FF34');
mysql> SELECT HEX(BIT_AND(a)), HEX(BIT_OR(a)), HEX(BIT_XOR(a))
       FROM t GROUP BY group_id;
+-----------------+----------------+-----------------+
| HEX(BIT_AND(a)) | HEX(BIT_OR(a)) | HEX(BIT_XOR(a)) |
+-----------------+----------------+-----------------+
| FFFFFFFFFFFF    | 000000000000   | 000000000000    |
| 1234            | FF34           | ED00            |
+-----------------+----------------+-----------------+
```

Special Handling of Hexadecimal Literals, Bit Literals, and NULL Literals

For backward compatibility, MySQL 8.0 evaluates bit operations in numeric context when all bit arguments are hexadecimal literals, bit literals, or `NULL` literals. That is, bit operations on binary-string

bit arguments do not use binary-string evaluation if all bit arguments are unadorned hexadecimal literals, bit literals, or `NULL` literals. (This does not apply to such literals if they are written with a `_binary` introducer, `BINARY` operator, or other way of specifying them explicitly as binary strings.)

The literal handling just described is the same as prior to MySQL 8.0. Examples:

- These bit operations evaluate the literals in numeric context and produce a `BIGINT` result:

```
b'0001' | b'0010'
X'0008' << 8
```

- These bit operations evaluate `NULL` in numeric context and produce a `BIGINT` result that has a `NULL` value:

```
NULL & NULL
NULL >> 4
```

In MySQL 8.0, you can cause those operations to evaluate the arguments in binary-string context by indicating explicitly that at least one argument is a binary string:

```
_binary b'0001' | b'0010'
_binary X'0008' << 8
BINARY NULL & NULL
BINARY NULL >> 4
```

The result of the last two expressions is `NULL`, just as without the `BINARY` operator, but the data type of the result is a binary string type rather than an integer type.

Bit-Operation Incompatibilities with MySQL 5.7

Because bit operations can handle binary string arguments natively in MySQL 8.0, some expressions produce a different result in MySQL 8.0 than in 5.7. The five problematic expression types to watch out for are:

```
nonliteral_binary { & | ^ } binary
binary  { & | ^ } nonliteral_binary
nonliteral_binary { << >> } anything
~ nonliteral_binary
AGGR_BIT_FUNC(nonliteral_binary)
```

Those expressions return `BIGINT` in MySQL 5.7, binary string in 8.0.

Explanation of notation:

- *`{ op1 op2 ... }`*: List of operators that apply to the given expression type.
- *`binary`*: Any kind of binary string argument, including a hexadecimal literal, bit literal, or `NULL` literal.
- *`nonliteral_binary`*: An argument that is a binary string value other than a hexadecimal literal, bit literal, or `NULL` literal.
- *`AGGR_BIT_FUNC`*: An aggregate function that takes bit-value arguments: `BIT_AND()`, `BIT_OR()`, `BIT_XOR()`.

For information about how to prepare in MySQL 5.7 for potential incompatibilities between MySQL 5.7 and 8.0, see Bit Functions and Operators, in MySQL 5.7 Reference Manual.

The following list describes available bit functions and operators:

- `|`

Bitwise OR.

The result type depends on whether the arguments are evaluated as binary strings or numbers:

- Binary-string evaluation occurs when the arguments have a binary string type, and at least one of them is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to unsigned 64-bit integers as necessary.
- Binary-string evaluation produces a binary string of the same length as the arguments. If the arguments have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

For more information, see the introductory discussion in this section.

```
mysql> SELECT 29 | 15;
        -> 31
mysql> SELECT _binary X'40404040' | X'01020304';
        -> 'ABCD'
```

- `&`

Bitwise AND.

The result type depends on whether the arguments are evaluated as binary strings or numbers:

- Binary-string evaluation occurs when the arguments have a binary string type, and at least one of them is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to unsigned 64-bit integers as necessary.
- Binary-string evaluation produces a binary string of the same length as the arguments. If the arguments have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

For more information, see the introductory discussion in this section.

```
mysql> SELECT 29 & 15;
        -> 13
mysql> SELECT HEX(_binary X'FF' & b'11110000');
        -> 'F0'
```

- `^`

Bitwise XOR.

The result type depends on whether the arguments are evaluated as binary strings or numbers:

- Binary-string evaluation occurs when the arguments have a binary string type, and at least one of them is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to unsigned 64-bit integers as necessary.
- Binary-string evaluation produces a binary string of the same length as the arguments. If the arguments have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

For more information, see the introductory discussion in this section.

```
mysql> SELECT 1 ^ 1;
        -> 0
mysql> SELECT 1 ^ 0;
        -> 1
mysql> SELECT 11 ^ 3;
```

```
        -> 8
mysql> SELECT HEX(_binary X'FEDC' ^ X'1111');
        -> 'EFCD'
```

- `<<`

  Shifts a longlong (`BIGINT`) number or binary string to the left.

  The result type depends on whether the bit argument is evaluated as a binary string or number:

  - Binary-string evaluation occurs when the bit argument has a binary string type, and is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to an unsigned 64-bit integer as necessary.
  - Binary-string evaluation produces a binary string of the same length as the bit argument. Numeric evaluation produces an unsigned 64-bit integer.

  Bits shifted off the end of the value are lost without warning, regardless of the argument type. In particular, if the shift count is greater or equal to the number of bits in the bit argument, all bits in the result are 0.

  For more information, see the introductory discussion in this section.

```
mysql> SELECT 1 << 2;
        -> 4
mysql> SELECT HEX(_binary X'00FF00FF00FF' << 8);
        -> 'FF00FF00FF00'
```

- `>>`

  Shifts a longlong (`BIGINT`) number or binary string to the right.

  The result type depends on whether the bit argument is evaluated as a binary string or number:

  - Binary-string evaluation occurs when the bit argument has a binary string type, and is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to an unsigned 64-bit integer as necessary.
  - Binary-string evaluation produces a binary string of the same length as the bit argument. Numeric evaluation produces an unsigned 64-bit integer.

  Bits shifted off the end of the value are lost without warning, regardless of the argument type. In particular, if the shift count is greater or equal to the number of bits in the bit argument, all bits in the result are 0.

  For more information, see the introductory discussion in this section.

```
mysql> SELECT 4 >> 2;
        -> 1
mysql> SELECT HEX(_binary X'00FF00FF00FF' >> 8);
        -> '0000FF00FF00'
```

- `~`

  Invert all bits.

  The result type depends on whether the bit argument is evaluated as a binary string or number:

  - Binary-string evaluation occurs when the bit argument has a binary string type, and is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument conversion to an unsigned 64-bit integer as necessary.

- Binary-string evaluation produces a binary string of the same length as the bit argument. Numeric evaluation produces an unsigned 64-bit integer.

For more information, see the introductory discussion in this section.

```
mysql> SELECT 5 & ~1;
        -> 4
mysql> SELECT HEX(~X'0000FFFF1111EEEE');
        -> 'FFFF0000EEEE1111'
```

- `BIT_COUNT(N)`

  Returns the number of bits that are set in the argument *`N`* as an unsigned 64-bit integer, or `NULL` if the argument is `NULL`.

```
mysql> SELECT BIT_COUNT(64), BIT_COUNT(BINARY 64);
        -> 1, 7
mysql> SELECT BIT_COUNT('64'), BIT_COUNT(_binary '64');
        -> 1, 7
mysql> SELECT BIT_COUNT(X'40'), BIT_COUNT(_binary X'40');
        -> 1, 1
```

12.13 Encryption and Compression Functions

Table 12.17 Encryption Functions

Name	Description
`AES_DECRYPT()`	Decrypt using AES
`AES_ENCRYPT()`	Encrypt using AES
`COMPRESS()`	Return result as a binary string
`DECODE()`	Decodes a string encrypted using ENCODE()
`DES_DECRYPT()`	Decrypt a string
`DES_ENCRYPT()`	Encrypt a string
`ENCODE()`	Encode a string
`ENCRYPT()`	Encrypt a string
`MD5()`	Calculate MD5 checksum
`PASSWORD()`	Calculate and return a password string
`RANDOM_BYTES()`	Return a random byte vector
`SHA1()`, `SHA()`	Calculate an SHA-1 160-bit checksum
`SHA2()`	Calculate an SHA-2 checksum
`UNCOMPRESS()`	Uncompress a string compressed
`UNCOMPRESSED_LENGTH()`	Return the length of a string before compression
`VALIDATE_PASSWORD_STRENGTH()`	Determine strength of password

Many encryption and compression functions return strings for which the result might contain arbitrary byte values. If you want to store these results, use a column with a `VARBINARY` or `BLOB` binary string data type. This will avoid potential problems with trailing space removal or character set conversion that would change data values, such as may occur if you use a nonbinary string data type (`CHAR`, `VARCHAR`, `TEXT`).

Some encryption functions return strings of ASCII characters: `MD5()`, `PASSWORD()`, `SHA()`, `SHA1()`, `SHA2()`. Their return value is a nonbinary string that has a character set and collation determined by the `character_set_connection` and `collation_connection` system variables.

For versions in which functions such as `MD5()` or `SHA1()` return a string of hexadecimal digits as a binary string, the return value cannot be converted to uppercase or compared in case-insensitive fashion as is. You must convert the value to a nonbinary string. See the discussion of binary string conversion in Section 12.10, “Cast Functions and Operators”.

If an application stores values from a function such as `MD5()` or `SHA1()` that returns a string of hex digits, more efficient storage and comparisons can be obtained by converting the hex representation to binary using `UNHEX()` and storing the result in a `BINARY(N)` column. Each pair of hexadecimal digits requires one byte in binary form, so the value of *`N`* depends on the length of the hex string. *`N`* is 16 for an `MD5()` value and 20 for a `SHA1()` value. For `SHA2()`, *`N`* ranges from 28 to 32 depending on the argument specifying the desired bit length of the result.

The size penalty for storing the hex string in a `CHAR` column is at least two times, up to eight times if the value is stored in a column that uses the `utf8` character set (where each character uses 4 bytes). Storing the string also results in slower comparisons because of the larger values and the need to take character set collation rules into account.

Suppose that an application stores `MD5()` string values in a `CHAR(32)` column:

```
CREATE TABLE md5_tbl (md5_val CHAR(32), ...);
INSERT INTO md5_tbl (md5_val, ...) VALUES(MD5('abcdef'), ...);
```

To convert hex strings to more compact form, modify the application to use `UNHEX()` and `BINARY(16)` instead as follows:

```
CREATE TABLE md5_tbl (md5_val BINARY(16), ...);
INSERT INTO md5_tbl (md5_val, ...) VALUES(UNHEX(MD5('abcdef')), ...);
```

Applications should be prepared to handle the very rare case that a hashing function produces the same value for two different input values. One way to make collisions detectable is to make the hash column a primary key.

Note

Exploits for the MD5 and SHA-1 algorithms have become known. You may wish to consider using another one-way encryption function described in this section instead, such as `SHA2()`.

Caution

Passwords or other sensitive values supplied as arguments to encryption functions are sent in cleartext to the MySQL server unless an SSL connection is used. Also, such values will appear in any MySQL logs to which they are written. To avoid these types of exposure, applications can encrypt sensitive values on the client side before sending them to the server. The same considerations apply to encryption keys. To avoid exposing these, applications can use stored procedures to encrypt and decrypt values on the server side.

- `AES_DECRYPT(`*`crypt_str`*`,`*`key_str`*`[,`*`init_vector`*`])`

  This function decrypts data using the official AES (Advanced Encryption Standard) algorithm. For more information, see the description of `AES_ENCRYPT()`.

  The optional initialization vector argument, *`init_vector`*. Statements that use `AES_DECRYPT()` are unsafe for statement-based replication and cannot be stored in the query cache.

- `AES_ENCRYPT(`*`str`*`,`*`key_str`*`[,`*`init_vector`*`])`

  `AES_ENCRYPT()` and `AES_DECRYPT()` implement encryption and decryption of data using the official AES (Advanced Encryption Standard) algorithm, previously known as “Rijndael.” The AES standard permits various key lengths. By default these functions implement AES with a 128-bit key

length. Key lengths of 196 or 256 bits can be used, as described later. The key length is a trade off between performance and security.

`AES_ENCRYPT()` encrypts the string *`str`* using the key string *`key_str`* and returns a binary string containing the encrypted output. `AES_DECRYPT()` decrypts the encrypted string *`crypt_str`* using the key string *`key_str`* and returns the original cleartext string. If either function argument is `NULL`, the function returns `NULL`.

The *`str`* and *`crypt_str`* arguments can be any length, and padding is automatically added to *`str`* so it is a multiple of a block as required by block-based algorithms such as AES. This padding is automatically removed by the `AES_DECRYPT()` function. The length of *`crypt_str`* can be calculated using this formula:

```
16 * (trunc(string_length / 16) + 1)
```

For a key length of 128 bits, the most secure way to pass a key to the *`key_str`* argument is to create a truly random 128-bit value and pass it as a binary value. For example:

```
INSERT INTO t
VALUES (1,AES_ENCRYPT('text',UNHEX('F3229A0B371ED2D9441B830D21A390C3')));
```

A passphrase can be used to generate an AES key by hashing the passphrase. For example:

```
INSERT INTO t
VALUES (1,AES_ENCRYPT('text', UNHEX(SHA2('My secret passphrase',512))));
```

Do not pass a password or passphrase directly to *`crypt_str`*, hash it first. Previous versions of this documentation suggested the former approach, but it is no longer recommended as the examples shown here are more secure.

If `AES_DECRYPT()` detects invalid data or incorrect padding, it returns `NULL`. However, it is possible for `AES_DECRYPT()` to return a non-`NULL` value (possibly garbage) if the input data or the key is invalid.

`AES_ENCRYPT()` and `AES_DECRYPT()` permit control of the block encryption mode and take an optional *`init_vector`* initialization vector argument:

- The `block_encryption_mode` system variable controls the mode for block-based encryption algorithms. Its default value is `aes-128-ecb`, which signifies encryption using a key length of 128 bits and ECB mode. For a description of the permitted values of this variable, see Section 5.1.5, “Server System Variables”.
- The optional *`init_vector`* argument provides an initialization vector for block encryption modes that require it.

For modes that require the optional *`init_vector`* argument, it must be 16 bytes or longer (bytes in excess of 16 are ignored). An error occurs if *`init_vector`* is missing.

For modes that do not require *`init_vector`*, it is ignored and a warning is generated if it is specified.

A random string of bytes to use for the initialization vector can be produced by calling `RANDOM_BYTES(16)`. For encryption modes that require an initialization vector, the same vector must be used for encryption and decryption.

```
mysql> SET block_encryption_mode = 'aes-256-cbc';
mysql> SET @key_str = SHA2('My secret passphrase',512);
mysql> SET @init_vector = RANDOM_BYTES(16);
mysql> SET @crypt_str = AES_ENCRYPT('text',@key_str,@init_vector);
```

```
mysql> SELECT AES_DECRYPT(@crypt_str,@key_str,@init_vector);
+-------------------------------------------------+
| AES_DECRYPT(@crypt_str,@key_str,@init_vector) |
+-------------------------------------------------+
| text                                            |
+-------------------------------------------------+
```

The following table lists each permitted block encryption mode, the SSL libraries that support it, and whether the initialization vector argument is required.

Block Encryption Mode	SSL Libraries that Support Mode	Initialization Vector Required
ECB	OpenSSL, yaSSL	No
CBC	OpenSSL, yaSSL	Yes
CFB1	OpenSSL	Yes
CFB8	OpenSSL	Yes
CFB128	OpenSSL	Yes
OFB	OpenSSL	Yes

Statements that use `AES_ENCRYPT()` or `AES_DECRYPT()` are unsafe for statement-based replication and cannot be stored in the query cache.

- `COMPRESS(`*`string_to_compress`*`)`

  Compresses a string and returns the result as a binary string. This function requires MySQL to have been compiled with a compression library such as `zlib`. Otherwise, the return value is always `NULL`. The compressed string can be uncompressed with `UNCOMPRESS()`.

```
mysql> SELECT LENGTH(COMPRESS(REPEAT('a',1000)));
        -> 21
mysql> SELECT LENGTH(COMPRESS(''));
        -> 0
mysql> SELECT LENGTH(COMPRESS('a'));
        -> 13
mysql> SELECT LENGTH(COMPRESS(REPEAT('a',16)));
        -> 15
```

  The compressed string contents are stored the following way:

  - Empty strings are stored as empty strings.
  - Nonempty strings are stored as a 4-byte length of the uncompressed string (low byte first), followed by the compressed string. If the string ends with space, an extra `.` character is added to avoid problems with endspace trimming should the result be stored in a `CHAR` or `VARCHAR` column. (However, use of nonbinary string data types such as `CHAR` or `VARCHAR` to store compressed strings is not recommended anyway because character set conversion may occur. Use a `VARBINARY` or `BLOB` binary string column instead.)

- `DECODE(`*`crypt_str`*`,`*`pass_str`*`)`

  `DECODE()` decrypts the encrypted string *`crypt_str`* using *`pass_str`* as the password. *`crypt_str`* should be a string returned from `ENCODE()`.

Note

The `ENCODE()` and `DECODE()` functions are deprecated, will be removed in a future MySQL release, and should no longer be used. Consider using `AES_ENCRYPT()` and `AES_DECRYPT()` instead.

- `DES_DECRYPT(`*`crypt_str`*`[,`*`key_str`*`])`

Decrypts a string encrypted with `DES_ENCRYPT()`. If an error occurs, this function returns `NULL`.

This function works only if MySQL has been configured with SSL support. See Section 6.4, “Using Secure Connections”.

If no *`key_str`* argument is given, `DES_DECRYPT()` examines the first byte of the encrypted string to determine the DES key number that was used to encrypt the original string, and then reads the key from the DES key file to decrypt the message. For this to work, the user must have the `SUPER` privilege. The key file can be specified with the `--des-key-file` server option.

If you pass this function a *`key_str`* argument, that string is used as the key for decrypting the message.

If the *`crypt_str`* argument does not appear to be an encrypted string, MySQL returns the given *`crypt_str`*.

Note

The `DES_ENCRYPT()` and `DES_DECRYPT()` functions are deprecated, will be removed in a future MySQL release, and should no longer be used. Consider using `AES_ENCRYPT()` and `AES_DECRYPT()` instead.

- `DES_ENCRYPT(`*`str`*`[,{`*`key_num`*`|`*`key_str`*`}])`

  Encrypts the string with the given key using the Triple-DES algorithm.

  This function works only if MySQL has been configured with SSL support. See Section 6.4, “Using Secure Connections”.

  The encryption key to use is chosen based on the second argument to `DES_ENCRYPT()`, if one was given. With no argument, the first key from the DES key file is used. With a *`key_num`* argument, the given key number (0 to 9) from the DES key file is used. With a *`key_str`* argument, the given key string is used to encrypt *`str`*.

  The key file can be specified with the `--des-key-file` server option.

  The return string is a binary string where the first character is `CHAR(128 | `*`key_num`*`)`. If an error occurs, `DES_ENCRYPT()` returns `NULL`.

  The 128 is added to make it easier to recognize an encrypted key. If you use a string key, *`key_num`* is 127.

  The string length for the result is given by this formula:

```
new_len = orig_len + (8 - (orig_len % 8)) + 1
```

  Each line in the DES key file has the following format:

```
key_num des_key_str
```

  Each *`key_num`* value must be a number in the range from `0` to `9`. Lines in the file may be in any order. *`des_key_str`* is the string that is used to encrypt the message. There should be at least one space between the number and the key. The first key is the default key that is used if you do not specify any key argument to `DES_ENCRYPT()`.

  You can tell MySQL to read new key values from the key file with the `FLUSH DES_KEY_FILE` statement. This requires the `RELOAD` privilege.

One benefit of having a set of default keys is that it gives applications a way to check for the existence of encrypted column values, without giving the end user the right to decrypt those values.

Note

The `DES_ENCRYPT()` and `DES_DECRYPT()` functions are deprecated, will be removed in a future MySQL release, and should no longer be used. Consider using `AES_ENCRYPT()` and `AES_DECRYPT()` instead.

```
mysql> SELECT customer_address FROM customer_table
     > WHERE crypted_credit_card = DES_ENCRYPT('credit_card_number');
```

- `ENCODE(str,pass_str)`

  `ENCODE()` encrypts *`str`* using *`pass_str`* as the password. The result is a binary string of the same length as *`str`*. To decrypt the result, use `DECODE()`.

  Note

  The `ENCODE()` and `DECODE()` functions are deprecated, will be removed in a future MySQL release, and should no longer be used.

  If you still need to use `ENCODE()`, a salt value must be used with it to reduce risk. For example:

  ```
  ENCODE('cleartext', CONCAT('my_random_salt','my_secret_password'))
  ```

  A new random salt value must be used whenever a password is updated.

- `ENCRYPT(str[,salt])`

  Encrypts *`str`* using the Unix `crypt()` system call and returns a binary string. The *`salt`* argument must be a string with at least two characters or the result will be `NULL`. If no *`salt`* argument is given, a random value is used.

  Note

  The `ENCRYPT()` function is deprecated, will be removed in a future MySQL release, and should no longer be used. Consider using `AES_ENCRYPT()` instead.

  ```
  mysql> SELECT ENCRYPT('hello');
         -> 'VxuFAJXVARROc'
  ```

  `ENCRYPT()` ignores all but the first eight characters of *`str`*, at least on some systems. This behavior is determined by the implementation of the underlying `crypt()` system call.

  The use of `ENCRYPT()` with the `ucs2`, `utf16`, `utf16le`, or `utf32` multibyte character sets is not recommended because the system call expects a string terminated by a zero byte.

  If `crypt()` is not available on your system (as is the case with Windows), `ENCRYPT()` always returns `NULL`.

- `MD5(str)`

  Calculates an MD5 128-bit checksum for the string. The value is returned as a string of 32 hexadecimal digits, or `NULL` if the argument was `NULL`. The return value can, for example, be used as a hash key. See the notes at the beginning of this section about storing hash values efficiently.

  The return value is a nonbinary string in the connection character set.

```
mysql> SELECT MD5('testing');
       -> 'ae2b1fca515949e5d54fb22b8ed95575'
```

This is the “RSA Data Security, Inc. MD5 Message-Digest Algorithm.”

See the note regarding the MD5 algorithm at the beginning this section.

- `PASSWORD(str)`

Note

This function is deprecated, will be removed in a future MySQL release, and should no longer be used.

Returns a hashed password string calculated from the cleartext password *`str`*. The return value is a nonbinary string in the connection character set, or `NULL` if the argument is `NULL`. This function is the SQL interface to the algorithm used by the server to encrypt MySQL passwords for storage in the `mysql.user` grant table.

The `old_passwords` system variable controls the password hashing method used by the `PASSWORD()` function. It also influences password hashing performed by `CREATE USER` and `GRANT` statements that specify a password using an `IDENTIFIED BY` clause.

The following table shows the permitted values of `old_passwords`, the password hashing method for each value, and which authentication plugins use passwords hashed with each method.

Value	Password Hashing Method	Associated Authentication Plugin
0	MySQL 4.1 native hashing	`mysql_native_password`
2	SHA-256 hashing	`sha256_password`

SHA-256 password hashing (`old_passwords=2`) uses a random salt value, which makes the result from `PASSWORD()` nondeterministic. Consequently, statements that use this function are not safe for statement-based replication and cannot be stored in the query cache.

Encryption performed by `PASSWORD()` is one-way (not reversible). It is not the same type of encryption used for Unix passwords; for that, use `ENCRYPT()`.

Note

`PASSWORD()` is used by the authentication system in MySQL Server; you should *not* use it in your own applications. For that purpose, consider a more secure function such as `AES_ENCRYPT()` or `SHA2()` instead. Also see RFC 2195, section 2 (Challenge-Response Authentication Mechanism (CRAM)), for more information about handling passwords and authentication securely in your applications.

Caution

Under some circumstances, statements that invoke `PASSWORD()` may be recorded in server logs or on the client side in a history file such as `~/.mysql_history`, which means that cleartext passwords may be read by anyone having read access to that information. For information about the conditions under which this occurs for the server logs and how to control it, see Section 6.1.2.3, “Passwords and Logging”. For similar information about client-side logging, see Section 4.5.1.3, “mysql Logging”.

- `RANDOM_BYTES(len)`

This function returns a binary string of *`len`* random bytes generated using the random number generator of the SSL library (OpenSSL or yaSSL). Permitted values of *`len`* range from 1 to 1024. For values outside that range, `RANDOM_BYTES()` generates a warning and returns `NULL`.

`RANDOM_BYTES()` can be used to provide the initialization vector for the `AES_DECRYPT()` and `AES_ENCRYPT()` functions. For use in that context, *`len`* must be at least 16. Larger values are permitted, but bytes in excess of 16 are ignored.

`RANDOM_BYTES()` generates a random value, which makes its result nondeterministic. Consequently, statements that use this function are unsafe for statement-based replication and cannot be stored in the query cache.

- `SHA1(str)`, `SHA(str)`

  Calculates an SHA-1 160-bit checksum for the string, as described in RFC 3174 (Secure Hash Algorithm). The value is returned as a string of 40 hexadecimal digits, or `NULL` if the argument was `NULL`. One of the possible uses for this function is as a hash key. See the notes at the beginning of this section about storing hash values efficiently. You can also use `SHA1()` as a cryptographic function for storing passwords. `SHA()` is synonymous with `SHA1()`.

  The return value is a nonbinary string in the connection character set.

```
mysql> SELECT SHA1('abc');
        -> 'a9993e364706816aba3e25717850c26c9cd0d89d'
```

  `SHA1()` can be considered a cryptographically more secure equivalent of `MD5()`. However, see the note regarding the MD5 and SHA-1 algorithms at the beginning this section.

- `SHA2(str, hash_length)`

  Calculates the SHA-2 family of hash functions (SHA-224, SHA-256, SHA-384, and SHA-512). The first argument is the cleartext string to be hashed. The second argument indicates the desired bit length of the result, which must have a value of 224, 256, 384, 512, or 0 (which is equivalent to 256). If either argument is `NULL` or the hash length is not one of the permitted values, the return value is `NULL`. Otherwise, the function result is a hash value containing the desired number of bits. See the notes at the beginning of this section about storing hash values efficiently.

  The return value is a nonbinary string in the connection character set.

```
mysql> SELECT SHA2('abc', 224);
        -> '23097d223405d8228642a477bda255b32aadbce4bda0b3f7e36c9da7'
```

  This function works only if MySQL has been configured with SSL support. See Section 6.4, “Using Secure Connections”.

  `SHA2()` can be considered cryptographically more secure than `MD5()` or `SHA1()`.

- `UNCOMPRESS(string_to_uncompress)`

  Uncompresses a string compressed by the `COMPRESS()` function. If the argument is not a compressed value, the result is `NULL`. This function requires MySQL to have been compiled with a compression library such as `zlib`. Otherwise, the return value is always `NULL`.

```
mysql> SELECT UNCOMPRESS(COMPRESS('any string'));
        -> 'any string'
mysql> SELECT UNCOMPRESS('any string');
        -> NULL
```

- `UNCOMPRESSED_LENGTH(compressed_string)`

Returns the length that the compressed string had before being compressed.

```
mysql> SELECT UNCOMPRESSED_LENGTH(COMPRESS(REPEAT('a',30)));
        -> 30
```

- `VALIDATE_PASSWORD_STRENGTH(str)`

  Given an argument representing a cleartext password, this function returns an integer to indicate how strong the password is. The return value ranges from 0 (weak) to 100 (strong).

  Password assessment by `VALIDATE_PASSWORD_STRENGTH()` is done by the `validate_password` plugin. If that plugin is not installed, the function always returns 0. For information about installing the `validate_password` plugin, see Section 6.5.3, “The Password Validation Plugin”. To examine or configure the parameters that affect password testing, check or set the system variables implemented by `validate_password` plugin. See Section 6.5.3.2, “Password Validation Plugin Options and Variables”.

  The password is subjected to increasingly strict tests and the return value reflects which tests were satisfied, as shown in the following table. In addition, if the `validate_password_check_user_name` system variable is enabled and the password matches the user name, `VALIDATE_PASSWORD_STRENGTH()` returns 0 regardless of how other `validate_password` system variables are set.

Password Test	Return Value
Length < 4	0
Length ≥ 4 and < `validate_password_length`	25
Satisfies policy 1 (`LOW`)	50
Satisfies policy 2 (`MEDIUM`)	75
Satisfies policy 3 (`STRONG`)	100

12.14 Information Functions

Table 12.18 Information Functions

Name	Description
`BENCHMARK()`	Repeatedly execute an expression
`CHARSET()`	Return the character set of the argument
`COERCIBILITY()`	Return the collation coercibility value of the string argument
`COLLATION()`	Return the collation of the string argument
`CONNECTION_ID()`	Return the connection ID (thread ID) for the connection
`CURRENT_ROLE()`	Returns the current active roles
`CURRENT_USER()`, `CURRENT_USER`	The authenticated user name and host name
`DATABASE()`	Return the default (current) database name
`FOUND_ROWS()`	For a SELECT with a LIMIT clause, the number of rows that would be returned were there no LIMIT clause
`LAST_INSERT_ID()`	Value of the AUTOINCREMENT column for the last INSERT
`ROLES_GRAPHML()`	Returns a GraphML document representing memory role subgraphs
`ROW_COUNT()`	The number of rows updated
`SCHEMA()`	Synonym for DATABASE()

Name	Description
SESSION_USER()	Synonym for USER()
SYSTEM_USER()	Synonym for USER()
USER()	The user name and host name provided by the client
VERSION()	Return a string that indicates the MySQL server version

- BENCHMARK(*count*,*expr*)

  The BENCHMARK() function executes the expression *expr* repeatedly *count* times. It may be used to time how quickly MySQL processes the expression. The result value is always 0. The intended use is from within the mysql client, which reports query execution times:

```
mysql> SELECT BENCHMARK(1000000,ENCODE('hello','goodbye'));
+--------------------------------------------+
| BENCHMARK(1000000,ENCODE('hello','goodbye')) |
+--------------------------------------------+
|                                          0 |
+--------------------------------------------+
1 row in set (4.74 sec)
```

  The time reported is elapsed time on the client end, not CPU time on the server end. It is advisable to execute BENCHMARK() several times, and to interpret the result with regard to how heavily loaded the server machine is.

  BENCHMARK() is intended for measuring the runtime performance of scalar expressions, which has some significant implications for the way that you use it and interpret the results:

  - Only scalar expressions can be used. Although the expression can be a subquery, it must return a single column and at most a single row. For example, BENCHMARK(10, (SELECT * FROM t)) will fail if the table t has more than one column or more than one row.
  - Executing a SELECT *expr* statement *N* times differs from executing SELECT BENCHMARK(*N*, *expr*) in terms of the amount of overhead involved. The two have very different execution profiles and you should not expect them to take the same amount of time. The former involves the parser, optimizer, table locking, and runtime evaluation *N* times each. The latter involves only runtime evaluation *N* times, and all the other components just once. Memory structures already allocated are reused, and runtime optimizations such as local caching of results already evaluated for aggregate functions can alter the results. Use of BENCHMARK() thus measures performance of the runtime component by giving more weight to that component and removing the “noise” introduced by the network, parser, optimizer, and so forth.

- CHARSET(*str*)

  Returns the character set of the string argument.

```
mysql> SELECT CHARSET('abc');
        -> 'utf8'
mysql> SELECT CHARSET(CONVERT('abc' USING latin1));
        -> 'latin1'
mysql> SELECT CHARSET(USER());
        -> 'utf8'
```

- COERCIBILITY(*str*)

  Returns the collation coercibility value of the string argument.

```
mysql> SELECT COERCIBILITY('abc' COLLATE utf8_swedish_ci);
        -> 0
mysql> SELECT COERCIBILITY(USER());
        -> 3
```

```
mysql> SELECT COERCIBILITY('abc');
        -> 4
mysql> SELECT COERCIBILITY(1000);
        -> 5
```

The return values have the meanings shown in the following table. Lower values have higher precedence.

Coercibility	Meaning	Example
0	Explicit collation	Value with COLLATE clause
1	No collation	Concatenation of strings with different collations
2	Implicit collation	Column value, stored routine parameter or local variable
3	System constant	USER() return value
4	Coercible	Literal string
5	Numeric	Numeric or temporal value
5	Ignorable	NULL or an expression derived from NULL

For more information, see Section 10.1.8.4, “Collation Coercibility in Expressions”.

- `COLLATION(str)`

  Returns the collation of the string argument.

```
mysql> SELECT COLLATION('abc');
        -> 'utf8_general_ci'
mysql> SELECT COLLATION(_utf8mb4'abc');
        -> 'utf8mb4_0900_ai_ci'
mysql> SELECT COLLATION(_latin1'abc');
        -> 'latin1_swedish_ci'
```

- `CONNECTION_ID()`

  Returns the connection ID (thread ID) for the connection. Every connection has an ID that is unique among the set of currently connected clients.

  The value returned by `CONNECTION_ID()` is the same type of value as displayed in the `ID` column of the `INFORMATION_SCHEMA.PROCESSLIST` table, the `Id` column of `SHOW PROCESSLIST` output, and the `PROCESSLIST_ID` column of the Performance Schema `threads` table.

```
mysql> SELECT CONNECTION_ID();
        -> 23786
```

- `CURRENT_ROLE()`

  Returns a `utf8` string containing the current active roles for the current session, separated by commas, or `NONE` if there are none. The value reflects the setting of the `sql_quote_show_create` system variable.

  Suppose that an account is granted roles as follows:

```
GRANT 'r1', 'r2' TO 'u1'@'localhost';
SET DEFAULT ROLE ALL TO 'u1'@'localhost';
```

  In sessions for `u1`, the initial `CURRENT_ROLE()` value is the default account role or roles. Using `SET ROLE` changes that:

```
mysql> SELECT CURRENT_ROLE();
+-------------------+
| CURRENT_ROLE()    |
+-------------------+
| `r1`@`%`,`r2`@`%` |
+-------------------+
mysql> SET ROLE 'r1'; SELECT CURRENT_ROLE();
+----------------+
| CURRENT_ROLE() |
+----------------+
| `r1`@`%`       |
+----------------+
```

- `CURRENT_USER`, `CURRENT_USER()`

  Returns the user name and host name combination for the MySQL account that the server used to authenticate the current client. This account determines your access privileges. The return value is a string in the `utf8` character set.

  The value of `CURRENT_USER()` can differ from the value of `USER()`.

```
mysql> SELECT USER();
        -> 'davida@localhost'
mysql> SELECT * FROM mysql.user;
ERROR 1044: Access denied for user ''@'localhost' to
database 'mysql'
mysql> SELECT CURRENT_USER();
        -> '@localhost'
```

  The example illustrates that although the client specified a user name of `davida` (as indicated by the value of the `USER()` function), the server authenticated the client using an anonymous user account (as seen by the empty user name part of the `CURRENT_USER()` value). One way this might occur is that there is no account listed in the grant tables for `davida`.

  Within a stored program or view, `CURRENT_USER()` returns the account for the user who defined the object (as given by its `DEFINER` value) unless defined with the `SQL SECURITY INVOKER` characteristic. In the latter case, `CURRENT_USER()` returns the object's invoker.

  Triggers and events have no option to define the `SQL SECURITY` characteristic, so for these objects, `CURRENT_USER()` returns the account for the user who defined the object. To return the invoker, use `USER()` or `SESSION_USER()`.

  The following statements support use of the `CURRENT_USER()` function to take the place of the name of (and, possibly, a host for) an affected user or a definer; in such cases, `CURRENT_USER()` is expanded where and as needed:

  - `DROP USER`
  - `RENAME USER`
  - `GRANT`
  - `REVOKE`
  - `CREATE FUNCTION`
  - `CREATE PROCEDURE`
  - `CREATE TRIGGER`
  - `CREATE EVENT`

- `CREATE VIEW`
- `ALTER EVENT`
- `ALTER VIEW`
- `SET PASSWORD`

For information about the implications that this expansion of `CURRENT_USER()` has for replication, see Section 18.4.1.8, “Replication of CURRENT_USER()”.

- `DATABASE()`

Returns the default (current) database name as a string in the `utf8` character set. If there is no default database, `DATABASE()` returns `NULL`. Within a stored routine, the default database is the database that the routine is associated with, which is not necessarily the same as the database that is the default in the calling context.

```
mysql> SELECT DATABASE();
        -> 'test'
```

If there is no default database, `DATABASE()` returns `NULL`.

- `FOUND_ROWS()`

A `SELECT` statement may include a `LIMIT` clause to restrict the number of rows the server returns to the client. In some cases, it is desirable to know how many rows the statement would have returned without the `LIMIT`, but without running the statement again. To obtain this row count, include an `SQL_CALC_FOUND_ROWS` option in the `SELECT` statement, and then invoke `FOUND_ROWS()` afterward:

```
mysql> SELECT SQL_CALC_FOUND_ROWS * FROM tbl_name
    -> WHERE id > 100 LIMIT 10;
mysql> SELECT FOUND_ROWS();
```

The second `SELECT` returns a number indicating how many rows the first `SELECT` would have returned had it been written without the `LIMIT` clause.

In the absence of the `SQL_CALC_FOUND_ROWS` option in the most recent successful `SELECT` statement, `FOUND_ROWS()` returns the number of rows in the result set returned by that statement. If the statement includes a `LIMIT` clause, `FOUND_ROWS()` returns the number of rows up to the limit. For example, `FOUND_ROWS()` returns 10 or 60, respectively, if the statement includes `LIMIT 10` or `LIMIT 50, 10`.

The row count available through `FOUND_ROWS()` is transient and not intended to be available past the statement following the `SELECT SQL_CALC_FOUND_ROWS` statement. If you need to refer to the value later, save it:

```
mysql> SELECT SQL_CALC_FOUND_ROWS * FROM ... ;
mysql> SET @rows = FOUND_ROWS();
```

If you are using `SELECT SQL_CALC_FOUND_ROWS`, MySQL must calculate how many rows are in the full result set. However, this is faster than running the query again without `LIMIT`, because the result set need not be sent to the client.

`SQL_CALC_FOUND_ROWS` and `FOUND_ROWS()` can be useful in situations when you want to restrict the number of rows that a query returns, but also determine the number of rows in the full result set without running the query again. An example is a Web script that presents a paged display

containing links to the pages that show other sections of a search result. Using FOUND_ROWS() enables you to determine how many other pages are needed for the rest of the result.

The use of SQL_CALC_FOUND_ROWS and FOUND_ROWS() is more complex for UNION statements than for simple SELECT statements, because LIMIT may occur at multiple places in a UNION. It may be applied to individual SELECT statements in the UNION, or global to the UNION result as a whole.

The intent of SQL_CALC_FOUND_ROWS for UNION is that it should return the row count that would be returned without a global LIMIT. The conditions for use of SQL_CALC_FOUND_ROWS with UNION are:

- The SQL_CALC_FOUND_ROWS keyword must appear in the first SELECT of the UNION.
- The value of FOUND_ROWS() is exact only if UNION ALL is used. If UNION without ALL is used, duplicate removal occurs and the value of FOUND_ROWS() is only approximate.
- If no LIMIT is present in the UNION, SQL_CALC_FOUND_ROWS is ignored and returns the number of rows in the temporary table that is created to process the UNION.

Beyond the cases described here, the behavior of FOUND_ROWS() is undefined (for example, its value following a SELECT statement that fails with an error).

Important

FOUND_ROWS() is not replicated reliably using statement-based replication. This function is automatically replicated using row-based replication.

- LAST_INSERT_ID(), LAST_INSERT_ID(*expr*)

With no argument, LAST_INSERT_ID() returns a BIGINT UNSIGNED (64-bit) value representing the first automatically generated value successfully inserted for an AUTO_INCREMENT column as a result of the most recently executed INSERT statement. The value of LAST_INSERT_ID() remains unchanged if no rows are successfully inserted.

With an argument, LAST_INSERT_ID() returns an unsigned integer.

For example, after inserting a row that generates an AUTO_INCREMENT value, you can get the value like this:

```
mysql> SELECT LAST_INSERT_ID();
        -> 195
```

The currently executing statement does not affect the value of LAST_INSERT_ID(). Suppose that you generate an AUTO_INCREMENT value with one statement, and then refer to LAST_INSERT_ID() in a multiple-row INSERT statement that inserts rows into a table with its own AUTO_INCREMENT column. The value of LAST_INSERT_ID() will remain stable in the second statement; its value for the second and later rows is not affected by the earlier row insertions. (However, if you mix references to LAST_INSERT_ID() and LAST_INSERT_ID(*expr*), the effect is undefined.)

If the previous statement returned an error, the value of LAST_INSERT_ID() is undefined. For transactional tables, if the statement is rolled back due to an error, the value of LAST_INSERT_ID() is left undefined. For manual ROLLBACK, the value of LAST_INSERT_ID() is not restored to that before the transaction; it remains as it was at the point of the ROLLBACK.

Within the body of a stored routine (procedure or function) or a trigger, the value of LAST_INSERT_ID() changes the same way as for statements executed outside the body of these kinds of objects. The effect of a stored routine or trigger upon the value of LAST_INSERT_ID() that is seen by following statements depends on the kind of routine:

- If a stored procedure executes statements that change the value of `LAST_INSERT_ID()`, the changed value is seen by statements that follow the procedure call.
- For stored functions and triggers that change the value, the value is restored when the function or trigger ends, so following statements will not see a changed value.

The ID that was generated is maintained in the server on a *per-connection basis*. This means that the value returned by the function to a given client is the first `AUTO_INCREMENT` value generated for most recent statement affecting an `AUTO_INCREMENT` column *by that client*. This value cannot be affected by other clients, even if they generate `AUTO_INCREMENT` values of their own. This behavior ensures that each client can retrieve its own ID without concern for the activity of other clients, and without the need for locks or transactions.

The value of `LAST_INSERT_ID()` is not changed if you set the `AUTO_INCREMENT` column of a row to a non-“magic” value (that is, a value that is not `NULL` and not `0`).

Important

If you insert multiple rows using a single `INSERT` statement, `LAST_INSERT_ID()` returns the value generated for the *first* inserted row *only*. The reason for this is to make it possible to reproduce easily the same `INSERT` statement against some other server.

For example:

```
mysql> USE test;
Database changed
mysql> CREATE TABLE t (
    ->   id INT AUTO_INCREMENT NOT NULL PRIMARY KEY,
    ->   name VARCHAR(10) NOT NULL
    -> );
Query OK, 0 rows affected (0.09 sec)

mysql> INSERT INTO t VALUES (NULL, 'Bob');
Query OK, 1 row affected (0.01 sec)

mysql> SELECT * FROM t;
+----+------+
| id | name |
+----+------+
|  1 | Bob  |
+----+------+
1 row in set (0.01 sec)

mysql> SELECT LAST_INSERT_ID();
+------------------+
| LAST_INSERT_ID() |
+------------------+
|                1 |
+------------------+
1 row in set (0.00 sec)

mysql> INSERT INTO t VALUES
    -> (NULL, 'Mary'), (NULL, 'Jane'), (NULL, 'Lisa');
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM t;
+----+------+
| id | name |
+----+------+
|  1 | Bob  |
|  2 | Mary |
|  3 | Jane |
|  4 | Lisa |
```

```
+----+------+
4 rows in set (0.01 sec)

mysql> SELECT LAST_INSERT_ID();
+------------------+
| LAST_INSERT_ID() |
+------------------+
|                2 |
+------------------+
1 row in set (0.00 sec)
```

Although the second `INSERT` statement inserted three new rows into `t`, the ID generated for the first of these rows was `2`, and it is this value that is returned by `LAST_INSERT_ID()` for the following `SELECT` statement.

If you use `INSERT IGNORE` and the row is ignored, the `LAST_INSERT_ID()` remains unchanged from the current value (or 0 is returned if the connection has not yet performed a successful `INSERT`) and, for non-transactional tables, the `AUTO_INCREMENT` counter is not incremented. For `InnoDB` tables, the `AUTO_INCREMENT` counter is incremented if `innodb_autoinc_lock_mode` is set to `1` or `2`, as demonstrated in the following example:

```
mysql> USE test;
Database changed

mysql> SELECT @@innodb_autoinc_lock_mode;
+----------------------------+
| @@innodb_autoinc_lock_mode |
+----------------------------+
|                          1 |
+----------------------------+
1 row in set (0.00 sec)

mysql> CREATE TABLE `t` (
`id` INT(11) NOT NULL AUTO_INCREMENT,
`val` INT(11) DEFAULT NULL,
PRIMARY KEY (`id`),
UNIQUE KEY `i1` (`val`)
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4;
Query OK, 0 rows affected (0.02 sec)

-- Insert two rows

mysql> INSERT INTO t (val) VALUES (1),(2);
Query OK, 2 rows affected (0.00 sec)
Records: 2  Duplicates: 0  Warnings: 0

-- With auto_increment_offset=1, the inserted rows
-- result in an AUTO_INCREMENT value of 3

mysql> SHOW CREATE TABLE t\G
*************************** 1. row ***************************
       Table: t
Create Table: CREATE TABLE `t` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `val` int(11) DEFAULT NULL,
  PRIMARY KEY (`id`),
  UNIQUE KEY `i1` (`val`)
) ENGINE=MyISAM AUTO_INCREMENT=3 DEFAULT CHARSET=utf8mb4
1 row in set (0.00 sec)

-- LAST_INSERT_ID() returns the first automatically generated
-- value that is successfully inserted for the AUTO_INCREMENT column

mysql> SELECT LAST_INSERT_ID();
+------------------+
| LAST_INSERT_ID() |
+------------------+
|                1 |
+------------------+
1 row in set (0.00 sec)
```

```
-- The attempted insertion of duplicate rows fail but errors are ignored

mysql> INSERT IGNORE INTO t (val) VALUES (1),(2);
Query OK, 0 rows affected (0.00 sec)
Records: 2  Duplicates: 2  Warnings: 0

-- With innodb_autoinc_lock_mode=1, the AUTO_INCREMENT counter
-- is incremented for the ignored rows

mysql> SHOW CREATE TABLE t\G
*************************** 1. row ***************************
       Table: t
Create Table: CREATE TABLE `t` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `val` int(11) DEFAULT NULL,
  PRIMARY KEY (`id`),
  UNIQUE KEY `i1` (`val`)
) ENGINE=MyISAM AUTO_INCREMENT=5 DEFAULT CHARSET=utf8mb4
1 row in set (0.00 sec)

-- The LAST_INSERT_ID is unchanged becuase the previous insert was unsuccessful

mysql> SELECT LAST_INSERT_ID();
+------------------+
| LAST_INSERT_ID() |
+------------------+
|                1 |
+------------------+
1 row in set (0.00 sec)
```

See Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB” for more information.

If *`expr`* is given as an argument to `LAST_INSERT_ID()`, the value of the argument is returned by the function and is remembered as the next value to be returned by `LAST_INSERT_ID()`. This can be used to simulate sequences:

1. Create a table to hold the sequence counter and initialize it:

   ```
   mysql> CREATE TABLE sequence (id INT NOT NULL);
   mysql> INSERT INTO sequence VALUES (0);
   ```

2. Use the table to generate sequence numbers like this:

   ```
   mysql> UPDATE sequence SET id=LAST_INSERT_ID(id+1);
   mysql> SELECT LAST_INSERT_ID();
   ```

   The `UPDATE` statement increments the sequence counter and causes the next call to `LAST_INSERT_ID()` to return the updated value. The `SELECT` statement retrieves that value. The `mysql_insert_id()` C API function can also be used to get the value. See Section 27.7.7.38, “mysql_insert_id()”.

You can generate sequences without calling `LAST_INSERT_ID()`, but the utility of using the function this way is that the ID value is maintained in the server as the last automatically generated value. It is multi-user safe because multiple clients can issue the `UPDATE` statement and get their own sequence value with the `SELECT` statement (or `mysql_insert_id()`), without affecting or being affected by other clients that generate their own sequence values.

Note that `mysql_insert_id()` is only updated after `INSERT` and `UPDATE` statements, so you cannot use the C API function to retrieve the value for `LAST_INSERT_ID(`*`expr`*`)` after executing other SQL statements like `SELECT` or `SET`.

- `ROLES_GRAPHML()`

Returns a `utf8` string containing a GraphML document representing memory role subgraphs. The `ROLE_ADMIN` or `SUPER` privilege is required to see content in the `<graphml>` element. Otherwise, the result shows only an empty element:

```
mysql> SELECT ROLES_GRAPHML();
+-------------------------------------------------+
| ROLES_GRAPHML()                                 |
+-------------------------------------------------+
| <?xml version="1.0" encoding="UTF-8"?><graphml /> |
+-------------------------------------------------+
```

- `ROW_COUNT()`

  `ROW_COUNT()` returns a value as follows:

  - DDL statements: 0. This applies to statements such as `CREATE TABLE` or `DROP TABLE`.
  - DML statements other than `SELECT`: The number of affected rows. This applies to statements such as `UPDATE`, `INSERT`, or `DELETE` (as before), but now also to statements such as `ALTER TABLE` and `LOAD DATA INFILE`.
  - `SELECT`: -1 if the statement returns a result set, or the number of rows “affected” if it does not. For example, for `SELECT * FROM t1`, `ROW_COUNT()` returns -1. For `SELECT * FROM t1 INTO OUTFILE 'file_name'`, `ROW_COUNT()` returns the number of rows written to the file.
  - `SIGNAL` statements: 0.

  For `UPDATE` statements, the affected-rows value by default is the number of rows actually changed. If you specify the `CLIENT_FOUND_ROWS` flag to `mysql_real_connect()` when connecting to `mysqld`, the affected-rows value is the number of rows “found”; that is, matched by the `WHERE` clause.

  For `REPLACE` statements, the affected-rows value is 2 if the new row replaced an old row, because in this case, one row was inserted after the duplicate was deleted.

  For `INSERT ... ON DUPLICATE KEY UPDATE` statements, the affected-rows value per row is 1 if the row is inserted as a new row, 2 if an existing row is updated, and 0 if an existing row is set to its current values. If you specify the `CLIENT_FOUND_ROWS` flag, the affected-rows value is 1 (not 0) if an existing row is set to its current values.

  The `ROW_COUNT()` value is similar to the value from the `mysql_affected_rows()` C API function and the row count that the `mysql` client displays following statement execution.

```
mysql> INSERT INTO t VALUES(1),(2),(3);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT ROW_COUNT();
+-------------+
| ROW_COUNT() |
+-------------+
|           3 |
+-------------+
1 row in set (0.00 sec)

mysql> DELETE FROM t WHERE i IN(1,2);
Query OK, 2 rows affected (0.00 sec)

mysql> SELECT ROW_COUNT();
+-------------+
| ROW_COUNT() |
+-------------+
|           2 |
```

```
+-------------+
1 row in set (0.00 sec)
```

Important

`ROW_COUNT()` is not replicated reliably using statement-based replication. This function is automatically replicated using row-based replication.

- `SCHEMA()`

  This function is a synonym for `DATABASE()`.

- `SESSION_USER()`

  `SESSION_USER()` is a synonym for `USER()`.

- `SYSTEM_USER()`

  `SYSTEM_USER()` is a synonym for `USER()`.

- `USER()`

  Returns the current MySQL user name and host name as a string in the `utf8` character set.

```
mysql> SELECT USER();
        -> 'davida@localhost'
```

  The value indicates the user name you specified when connecting to the server, and the client host from which you connected. The value can be different from that of `CURRENT_USER()`.

- `VERSION()`

  Returns a string that indicates the MySQL server version. The string uses the `utf8` character set. The value might have a suffix in addition to the version number. See the description of the `version` system variable in Section 5.1.5, “Server System Variables”.

  This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

```
mysql> SELECT VERSION();
        -> '8.0.3-standard'
```

12.15 Spatial Analysis Functions

MySQL provides functions to perform various operations on spatial data. These functions can be grouped into several major categories according to the type of operation they perform:

- Functions that create geometries in various formats (WKT, WKB, internal)
- Functions that convert geometries between formats
- Functions that access qualitative or quantitative properties of a geometry
- Functions that describe relations between two geometries
- Functions that create new geometries from existing ones

For general background about MySQL support for using spatial data, see Section 11.5, “Extensions for Spatial Data”.

12.15.1 Spatial Function Reference

The following table lists each spatial function and provides a short description of each one.

Table 12.19 Spatial Functions

Name	Description
GeometryCollection()	Construct geometry collection from geometries
LineString()	Construct LineString from Point values
MBRContains()	Whether MBR of one geometry contains MBR of another
MBRCoveredBy()	Whether one MBR is covered by another
MBRCovers()	Whether one MBR covers another
MBRDisjoint()	Whether MBRs of two geometries are disjoint
MBREquals()	Whether MBRs of two geometries are equal
MBRIntersects()	Whether MBRs of two geometries intersect
MBROverlaps()	Whether MBRs of two geometries overlap
MBRTouches()	Whether MBRs of two geometries touch
MBRWithin()	Whether MBR of one geometry is within MBR of another
MultiLineString()	Contruct MultiLineString from LineString values
MultiPoint()	Construct MultiPoint from Point values
MultiPolygon()	Construct MultiPolygon from Polygon values
Point()	Construct Point from coordinates
Polygon()	Construct Polygon from LineString arguments
ST_Area()	Return Polygon or MultiPolygon area
ST_AsBinary(), ST_AsWKB()	Convert from internal geometry format to WKB
ST_AsGeoJSON()	Generate GeoJSON object from geometry
ST_AsText(), ST_AsWKT()	Convert from internal geometry format to WKT
ST_Buffer()	Return geometry of points within given distance from geometry
ST_Buffer_Strategy()	Produce strategy option for ST_Buffer()
ST_Centroid()	Return centroid as a point
ST_Contains()	Whether one geometry contains another
ST_ConvexHull()	Return convex hull of geometry
ST_Crosses()	Whether one geometry crosses another
ST_Difference()	Return point set difference of two geometries
ST_Dimension()	Dimension of geometry
ST_Disjoint()	Whether one geometry is disjoint from another
ST_Distance()	The distance of one geometry from another
ST_Distance_Sphere()	Minimum distance on earth between two geometries
ST_EndPoint()	End Point of LineString
ST_Envelope()	Return MBR of geometry
ST_Equals()	Whether one geometry is equal to another
ST_ExteriorRing()	Return exterior ring of Polygon
ST_GeoHash()	Produce a geohash value
ST_GeomCollFromText(), ST_GeometryCollectionFromText(), ST_GeomCollFromTxt()	Return geometry collection from WKT

Name	Description
ST_GeomCollFromWKB(), ST_GeometryCollectionFromWKB()	Return geometry collection from WKB
ST_GeometryN()	Return N-th geometry from geometry collection
ST_GeometryType()	Return name of geometry type
ST_GeomFromGeoJSON()	Generate geometry from GeoJSON object
ST_GeomFromText(), ST_GeometryFromText()	Return geometry from WKT
ST_GeomFromWKB(), ST_GeometryFromWKB()	Return geometry from WKB
ST_InteriorRingN()	Return N-th interior ring of Polygon
ST_Intersection()	Return point set intersection of two geometries
ST_Intersects()	Whether one geometry intersects another
ST_IsClosed()	Whether a geometry is closed and simple
ST_IsEmpty()	Placeholder function
ST_IsSimple()	Whether a geometry is simple
ST_IsValid()	Whether a geometry is valid
ST_LatFromGeoHash()	Return latitude from geohash value
ST_Length()	Return length of LineString
ST_LineFromText(), ST_LineStringFromText()	Construct LineString from WKT
ST_LineFromWKB(), ST_LineStringFromWKB()	Construct LineString from WKB
ST_LongFromGeoHash()	Return longitude from geohash value
ST_MakeEnvelope()	Rectangle around two points
ST_MLineFromText(), ST_MultiLineStringFromText()	Construct MultiLineString from WKT
ST_MLineFromWKB(), ST_MultiLineStringFromWKB()	Construct MultiLineString from WKB
ST_MPointFromText(), ST_MultiPointFromText()	Construct MultiPoint from WKT
ST_MPointFromWKB(), ST_MultiPointFromWKB()	Construct MultiPoint from WKB
ST_MPolyFromText(), ST_MultiPolygonFromText()	Construct MultiPolygon from WKT
ST_MPolyFromWKB(), ST_MultiPolygonFromWKB()	Construct MultiPolygon from WKB
ST_NumGeometries()	Return number of geometries in geometry collection
ST_NumInteriorRing(), ST_NumInteriorRings()	Return number of interior rings in Polygon
ST_NumPoints()	Return number of points in LineString
ST_Overlaps()	Whether one geometry overlaps another
ST_PointFromGeoHash()	Convert geohash value to POINT value
ST_PointFromText()	Construct Point from WKT
ST_PointFromWKB()	Construct Point from WKB

Name	Description
ST_PointN()	Return N-th point from LineString
ST_PolyFromText(), ST_PolygonFromText()	Construct Polygon from WKT
ST_PolyFromWKB(), ST_PolygonFromWKB()	Construct Polygon from WKB
ST_Simplify()	Return simplified geometry
ST_SRID()	Return spatial reference system ID for geometry
ST_StartPoint()	Start Point of LineString
ST_SwapXY()	Return argument with X/Y coordinates swapped
ST_SymDifference()	Return point set symmetric difference of two geometries
ST_Touches()	Whether one geometry touches another
ST_Union()	Return point set union of two geometries
ST_Validate()	Return validated geometry
ST_Within()	Whether one geometry is within another
ST_X()	Return X coordinate of Point
ST_Y()	Return Y coordinate of Point

12.15.2 Argument Handling by Spatial Functions

Spatial values, or geometries, have the properties described at Section 11.5.2.2, “Geometry Class”. The following discussion lists general spatial function argument-handling characteristics. Specific functions or groups of functions may have additional argument-handling characteristics, as discussed in the sections where those function descriptions occur.

Spatial functions are defined only for valid geometry values.

The spatial reference identifier (SRID) of a geometry identifies the coordinate space in which the geometry is defined. In MySQL, the SRID value is an integer associated with the geometry value. The maximum usable SRID value is $2^{32}-1$. If a larger value is given, only the lower 32 bits are used.

In MySQL, all computations are done assuming SRID 0, regardless of the actual SRID value. SRID 0 represents an infinite flat Cartesian plane with no units assigned to its axes. In the future, computations may use the specified SRID values. To ensure SRID 0 behavior, create geometry values using SRID 0. SRID 0 is the default for new geometry values if no SRID is specified.

Geometry values produced by any spatial function inherit the SRID of the geometry arguments.

Spatial functions that take multiple geometry arguments require those arguments to have the same SRID value (that is, same value in the lower 32 bits). Assuming equal SRIDs, spatial functions do nothing with them after performing the equality check; geometry values are implicitly handled using Cartesian coordinates (SRID 0). If a spatial function returns ER_GIS_DIFFERENT_SRIDS, it means that the geometry arguments did not all have the same SRID. You must modify them to have the same SRID.

The Open Geospatial Consortium guidelines require that input polygons already be closed, so unclosed polygons are rejected as invalid rather than being closed.

Empty geometry-collection handling is as follows: An empty WKT input geometry collection may be specified as 'GEOMETRYCOLLECTION()'. This is also the output WKT resulting from a spatial operation that produces an empty geometry collection.

During parsing of a nested geometry collection, the collection is flattened and its basic components are used in various GIS operations to compute results. This provides additional flexibility to users because

it is unnecessary to be concerned about the uniqueness of geometry data. Nested geometry collections may be produced from nested GIS function calls without having to be explicitly flattened first.

12.15.3 Functions That Create Geometry Values from WKT Values

These functions take as arguments a Well-Known Text (WKT) representation and, optionally, a spatial reference system identifier (SRID). They return the corresponding geometry.

The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

`ST_GeomFromText()` accepts a WKT value of any geometry type as its first argument. Other functions provide type-specific construction functions for construction of geometry values of each geometry type.

Functions such as `ST_GeomFromText()` that accept WKT geometry collection arguments understand both OpenGIS `'GEOMETRYCOLLECTION EMPTY'` standard syntax and MySQL `'GEOMETRYCOLLECTION()'` nonstandard syntax. Functions such as `ST_AsWKT()` that produce WKT values produce `'GEOMETRYCOLLECTION EMPTY'` standard syntax:

```
mysql> SET @s1 = ST_GeomFromText('GEOMETRYCOLLECTION()');
mysql> SET @s2 = ST_GeomFromText('GEOMETRYCOLLECTION EMPTY');
mysql> SELECT ST_AsWKT(@s1), ST_AsWKT(@s2);
+--------------------------+--------------------------+
| ST_AsWKT(@s1)            | ST_AsWKT(@s2)            |
+--------------------------+--------------------------+
| GEOMETRYCOLLECTION EMPTY | GEOMETRYCOLLECTION EMPTY |
+--------------------------+--------------------------+
```

For a description of WKT format, see Well-Known Text (WKT) Format.

- `ST_GeomCollFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_GeometryCollectionFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_GeomCollFromTxt(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`

  Constructs a `GeometryCollection` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

  ```
  mysql> SET @g = "MULTILINESTRING((10 10, 11 11), (9 9, 10 10))";
  mysql> SELECT ST_AsText(ST_GeomCollFromText(@g));
  +-----------------------------------------+
  | ST_AsText(ST_GeomCollFromText(@g))      |
  +-----------------------------------------+
  | MULTILINESTRING((10 10,11 11),(9 9,10 10)) |
  +-----------------------------------------+
  ```

- `ST_GeomFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_GeometryFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`

  Constructs a geometry value of any type using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_LineFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_LineStringFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`

  Constructs a `LineString` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_MLineFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_MultiLineStringFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`

  Constructs a `MultiLineString` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_MPointFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`, `ST_MultiPointFromText(`*`wkt`*`[, `*`srid`*` [, `*`options`*`]])`

  Constructs a `MultiPoint` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

  In MySQL 8.0, spatial functions such as `ST_MPointFromText()` and `ST_GeomFromText()` that accept WKT-format representations of `MultiPoint` values permit individual points within values to be surrounded by parentheses. For example, both of the following function calls are valid:

```
ST_MPointFromText('MULTIPOINT (1 1, 2 2, 3 3)')
ST_MPointFromText('MULTIPOINT ((1 1), (2 2), (3 3))')
```

- `ST_MPolyFromText(wkt[, srid [, options]])`, `ST_MultiPolygonFromText(wkt[, srid [, options]])`

  Constructs a `MultiPolygon` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_PointFromText(wkt[, srid [, options]])`

  Constructs a `Point` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_PolyFromText(wkt[, srid [, options]])`, `ST_PolygonFromText(wkt[, srid [, options]])`

  Constructs a `Polygon` value using its WKT representation and SRID.

  The result is `NULL` if the geometry argument is `NULL` or not a syntactically well-formed geometry, or if the SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

12.15.4 Functions That Create Geometry Values from WKB Values

These functions take as arguments a `BLOB` containing a Well-Known Binary (WKB) representation and, optionally, a spatial reference system identifier (SRID). They return the corresponding geometry.

The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

`ST_GeomFromWKB()` accepts a WKB value of any geometry type as its first argument. Other functions provide type-specific construction functions for construction of geometry values of each geometry type.

Prior to MySQL 8.0, these functions also accepted geometry objects as returned by the functions in Section 12.15.5, “MySQL-Specific Functions That Create Geometry Values”. Geometry arguments are no longer permitted and produce an error. To migrate calls from using geometry arguments to using WKB arguments, follow these guidelines:

- Rewrite constructs such as `ST_GeomFromWKB(Point(0, 0))` as `Point(0, 0)`.

- Rewrite constructs such as `ST_GeomFromWKB(Point(0, 0), 4326)` as `ST_SRID(Point(0, 0), 4326)` or `ST_GeomFromWKB(ST_AsWKB(Point(0, 0)), 4326)`.

For a description of WKB format, see Well-Known Binary (WKB) Format.

- `ST_GeomCollFromWKB(wkb[, srid [, options]])`, `ST_GeometryCollectionFromWKB(wkb[, srid [, options]])`

  Constructs a `GeometryCollection` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_GeomFromWKB(wkb[, srid [, options]])`, `ST_GeometryFromWKB(wkb[, srid [, options]])`

  Constructs a geometry value of any type using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_LineFromWKB(wkb[, srid [, options]])`, `ST_LineStringFromWKB(wkb[, srid [, options]])`

  Constructs a `LineString` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_MLineFromWKB(wkb[, srid [, options]])`, `ST_MultiLineStringFromWKB(wkb[, srid [, options]])`

  Constructs a `MultiLineString` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_MPointFromWKB(wkb[, srid [, options]])`, `ST_MultiPointFromWKB(wkb[, srid [, options]])`

  Constructs a `MultiPoint` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_MPolyFromWKB(wkb[, srid [, options]])`, `ST_MultiPolygonFromWKB(wkb[, srid [, options]])`

  Constructs a `MultiPolygon` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_PointFromWKB(wkb[, srid [, options]])`

  Constructs a `Point` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

- `ST_PolyFromWKB(wkb[, srid [, options]])`, `ST_PolygonFromWKB(wkb[, srid [, options]])`

  Constructs a `Polygon` value using its WKB representation and SRID.

  The result is `NULL` if the WKB or SRID argument is `NULL`.

  The functions interpret geographic coordinates (latitude, longitude) as in the order specified by the spatial reference system. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

12.15.5 MySQL-Specific Functions That Create Geometry Values

MySQL provides a set of useful nonstandard functions for creating geometry values. The functions described in this section are MySQL extensions to the OpenGIS specification.

These functions produce geometry objects from either WKB values or geometry objects as arguments. If any argument is not a proper WKB or geometry representation of the proper object type, the return value is `NULL`.

For example, you can insert the geometry return value from `Point()` directly into a `POINT` column:

```
INSERT INTO t1 (pt_col) VALUES(Point(1,2));
```

- `GeometryCollection(g1, g2, ...)`

  Constructs a `GeometryCollection`.

  `GeometryCollection()` returns all the proper geometries contained in the argument even if a nonsupported geometry is present.

  `GeometryCollection()` with no arguments is permitted as a way to create an empty geometry.

- `LineString(pt1, pt2, ...)`

  Constructs a `LineString` value from a number of `Point` or WKB `Point` arguments. If the number of arguments is less than two, the return value is `NULL`.

- `MultiLineString(ls1, ls2, ...)`

  Constructs a `MultiLineString` value using `LineString` or WKB `LineString` arguments.

- `MultiPoint(pt1, pt2, ...)`

  Constructs a `MultiPoint` value using `Point` or WKB `Point` arguments.

- `MultiPolygon(poly1, poly2, ...)`

  Constructs a `MultiPolygon` value from a set of `Polygon` or WKB `Polygon` arguments.

- `Point(x, y)`

  Constructs a `Point` using its coordinates.

- `Polygon(ls1, ls2, ...)`

  Constructs a `Polygon` value from a number of `LineString` or WKB `LineString` arguments. If any argument does not represent a `LinearRing` (that is, not a closed and simple `LineString`), the return value is `NULL`.

12.15.6 Geometry Format Conversion Functions

MySQL supports the functions listed in this section for converting geometry values from internal geometry format to WKT or WKB format, or for swapping the order of X and Y coordinates.

There are also functions to convert a string from WKT or WKB format to internal geometry format. See Section 12.15.3, “Functions That Create Geometry Values from WKT Values”, and Section 12.15.4, “Functions That Create Geometry Values from WKB Values”.

Functions such as `ST_GeomFromText()` that accept WKT geometry collection arguments understand both OpenGIS `'GEOMETRYCOLLECTION EMPTY'` standard syntax and MySQL `'GEOMETRYCOLLECTION()'` nonstandard syntax. Functions such as `ST_AsWKT()` that produce WKT values produce `'GEOMETRYCOLLECTION EMPTY'` standard syntax:

```
mysql> SET @s1 = ST_GeomFromText('GEOMETRYCOLLECTION()');
mysql> SET @s2 = ST_GeomFromText('GEOMETRYCOLLECTION EMPTY');
mysql> SELECT ST_AsWKT(@s1), ST_AsWKT(@s2);
+-------------------------+-------------------------+
| ST_AsWKT(@s1)           | ST_AsWKT(@s2)           |
```

```
+-------------------------+-------------------------+
| GEOMETRYCOLLECTION EMPTY | GEOMETRYCOLLECTION EMPTY |
+-------------------------+-------------------------+
```

- `ST_AsBinary(g [, options])`, `ST_AsWKB(g [, options])`

  Converts a value in internal geometry format to its WKB representation and returns the binary result.

  The result is `NULL` if the geometry argument is `NULL`. If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  The function return value has geographic coordinates (latitude, longitude) in the order specified by the spatial reference system that applies to the geometry argument. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

```
mysql> SET @g = ST_LineFromText('LINESTRING(0 5,5 10,10 15)', 4326);
mysql> SELECT ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g)));
+----------------------------------------+
| ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g))) |
+----------------------------------------+
| LINESTRING(5 0,10 5,15 10)             |
+----------------------------------------+
mysql> SELECT ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g, 'axis-order=long-lat')));
+-------------------------------------------------------------+
| ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g, 'axis-order=long-lat'))) |
+-------------------------------------------------------------+
| LINESTRING(0 5,5 10,10 15)                                  |
+-------------------------------------------------------------+
mysql> SELECT ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g, 'axis-order=lat-long')));
+-------------------------------------------------------------+
| ST_AsText(ST_GeomFromWKB(ST_AsWKB(@g, 'axis-order=lat-long'))) |
+-------------------------------------------------------------+
| LINESTRING(5 0,10 5,15 10)                                  |
+-------------------------------------------------------------+
```

- `ST_AsText(g [, options])`, `ST_AsWKT(g [, options])`

  Converts a value in internal geometry format to its WKT representation and returns the string result.

  The result is `NULL` if the geometry argument is `NULL`. If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  The function return value has geographic coordinates (latitude, longitude) in the order specified by the spatial reference system that applies to the geometry argument. An optional *`options`* argument may be given to override the default axis order.

  The result is `NULL` if the *`options`* argument is `NULL`. If the *`options`* argument is invalid, an error occurs to indicate why.

```
mysql> SET @g = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_AsText(ST_GeomFromText(@g));
+------------------------------+
| ST_AsText(ST_GeomFromText(@g)) |
+------------------------------+
| LINESTRING(1 1,2 2,3 3)      |
+------------------------------+
```

  Output for `MultiPoint` values includes parentheses around each point. For example:

```
mysql> SELECT ST_AsText(ST_GeomFromText(@mp));
```

```
+--------------------------------+
| ST_AsText(ST_GeomFromText(@mp)) |
+--------------------------------+
| MULTIPOINT((1 1),(2 2),(3 3))   |
+--------------------------------+
```

- `ST_SwapXY(g)`

  Accepts an argument in internal geometry format, swaps the X and Y values of each coordinate pair within the geometry, and returns the result.

  The result is `NULL` if the argument is `NULL`. If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

```
mysql> SET @g = ST_LineFromText('LINESTRING(0 5,5 10,10 15)');
mysql> SELECT ST_AsText(@g);
+---------------------------+
| ST_AsText(@g)             |
+---------------------------+
| LINESTRING(0 5,5 10,10 15) |
+---------------------------+
mysql> SELECT ST_AsText(ST_SwapXY(@g));
+---------------------------+
| ST_AsText(ST_SwapXY(@g))   |
+---------------------------+
| LINESTRING(5 0,10 5,15 10) |
+---------------------------+
```

12.15.7 Geometry Property Functions

Each function that belongs to this group takes a geometry value as its argument and returns some quantitative or qualitative property of the geometry. Some functions restrict their argument type. Such functions return `NULL` if the argument is of an incorrect geometry type. For example, the `ST_Area()` polygon function returns `NULL` if the object type is neither `Polygon` nor `MultiPolygon`.

12.15.7.1 General Geometry Property Functions

The functions listed in this section do not restrict their argument and accept a geometry value of any type.

- `ST_Dimension(g)`

  Returns the inherent dimension of the geometry value *g*, or `NULL` if the argument is `NULL`. The dimension can be −1, 0, 1, or 2. The meaning of these values is given in Section 11.5.2.2, “Geometry Class”.

```
mysql> SELECT ST_Dimension(ST_GeomFromText('LineString(1 1,2 2)'));
+------------------------------------------------------+
| ST_Dimension(ST_GeomFromText('LineString(1 1,2 2)')) |
+------------------------------------------------------+
|                                                    1 |
+------------------------------------------------------+
```

- `ST_Envelope(g)`

  Returns the minimum bounding rectangle (MBR) for the geometry value *g*, or `NULL` if the argument is `NULL`. The result is returned as a `Polygon` value that is defined by the corner points of the bounding box:

```
POLYGON((MINX MINY, MAXX MINY, MAXX MAXY, MINX MAXY, MINX MINY))
```

```
mysql> SELECT ST_AsText(ST_Envelope(ST_GeomFromText('LineString(1 1,2 2)')));
+-------------------------------------------------------------------+
```

```
| ST_AsText(ST_Envelope(ST_GeomFromText('LineString(1 1,2 2)'))) |
+----------------------------------------------------------------+
| POLYGON((1 1,2 1,2 2,1 2,1 1))                                 |
+----------------------------------------------------------------+
```

If the argument is a point or a vertical or horizontal line segment, `ST_Envelope()` returns the point or the line segment as its MBR rather than returning an invalid polygon:

```
mysql> SELECT ST_AsText(ST_Envelope(ST_GeomFromText('LineString(1 1,1 2)')));
+----------------------------------------------------------------+
| ST_AsText(ST_Envelope(ST_GeomFromText('LineString(1 1,1 2)'))) |
+----------------------------------------------------------------+
| LINESTRING(1 1,1 2)                                            |
+----------------------------------------------------------------+
```

- `ST_GeometryType(g)`

  Returns a binary string indicating the name of the geometry type of which the geometry instance *g* is a member, or `NULL` if the argument is `NULL`. The name corresponds to one of the instantiable `Geometry` subclasses.

```
mysql> SELECT ST_GeometryType(ST_GeomFromText('POINT(1 1)'));
+-----------------------------------------------+
| ST_GeometryType(ST_GeomFromText('POINT(1 1)')) |
+-----------------------------------------------+
| POINT                                         |
+-----------------------------------------------+
```

- `ST_IsEmpty(g)`

  This function is a placeholder that returns 0 for any valid geometry value, 1 for any invalid geometry value, or `NULL` if the argument is `NULL`.

  MySQL does not support GIS `EMPTY` values such as `POINT EMPTY`.

- `ST_IsSimple(g)`

  Returns 1 if the geometry value *g* has no anomalous geometric points, such as self-intersection or self-tangency. `ST_IsSimple()` returns 0 if the argument is not simple, and `NULL` if the argument is `NULL`.

  The description of each instantiable geometric class given earlier in the chapter includes the specific conditions that cause an instance of that class to be classified as not simple. (See Section 11.5.2.1, “The Geometry Class Hierarchy”.)

- `ST_SRID(g[, new_srid_val])`

  With a single argument representing a valid geometry object *g*, `ST_SRID()` returns an integer indicating the spatial reference system ID associated with *g*.

  The result is `NULL` if the geometry argument is `NULL`. If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  If the optional second argument representing a valid SRID value is given, `ST_SRID()` returns an object with the same type as its first argument, having an SRID value equal to the second argument, or `NULL` if either argument is `NULL`. These conditions apply when the second argument is given:

  - Error handling for the first argument is as for the single-argument syntax, as described previously.
  - If the SRID value in the second argument is not within the range of a 32-bit unsigned integer, an `ER_DATA_OUT_OF_RANGE` error occurs.

- If the SRID value in the second argument does not refer to a known spatial reference system, an `ER_SRS_NOT_FOUND` error occurs.

```
mysql> SET @g = ST_GeomFromText('LineString(1 1,2 2)', 0);
mysql> SELECT ST_SRID(@g);
+-------------+
| ST_SRID(@g) |
+-------------+
|           0 |
+-------------+
mysql> SET @g2 = ST_SRID(@g, 4326);
mysql> SELECT ST_SRID(@g2);
+--------------+
| ST_SRID(@g2) |
+--------------+
|         4326 |
+--------------+
```

It is possible to create a geometry in a particular SRID by passing to `ST_SRID()` the result of one of the MySQL-specific functions for creating spatial values, along with an SRID value. For example:

```
SET @g1 = ST_SRID(Point(1, 1), 4326);
```

However, that method creates the geometry in SRID 0, then casts it to SRID 4326 (WGS 84). A preferable alternative is to create the geometry with the correct spatial reference system (SRS) to begin with. For example:

```
SET @g1 = ST_PointFromText('POINT(1 1)', 4326);
SET @g1 = ST_GeomFromText('POINT(1 1)', 4326);
```

The two-argument form of `ST_SRID()` is useful for tasks such as correcting or changing the SRS of geometries that have an incorrect SRID.

12.15.7.2 Point Property Functions

A `Point` consists of X and Y coordinates, which may be obtained using the `ST_X()` and `ST_Y()` functions, respectively. These functions also permit an optional second argument that specifies an X or Y coordinate value, in which case the function result is the `Point` object from the first argument with the appropriate coordinate modified to be equal to the second argument.

- `ST_X(`*p*`[, `*new_x_val*`])`

  With a single argument representing a valid `Point` object *p*, `ST_X()` returns the X-coordinate value of *p* as a double-precision number, or `NULL` if the argument is `NULL`. These conditions apply:

  - If the argument is a valid geometry but not a `Point` object, an `ER_UNEXPECTED_GEOMETRY_TYPE` error occurs.
  - If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  If the optional second argument representing a valid X-coordinate value is given, `ST_X()` returns a `Point` object with an X coordinate equal to the second argument, or `NULL` if either argument is `NULL`. These conditions apply:

  - Error handling for the first argument is as for the single-argument syntax, as described previously.
  - If the coordinate value in the second argument is `-inf`, `+inf`, or `NaN`, an `ER_DATA_OUT_OF_RANGE` error occurs.

```
mysql> SELECT ST_X(Point(56.7, 53.34));
+--------------------------+
| ST_X(Point(56.7, 53.34)) |
+--------------------------+
|                     56.7 |
+--------------------------+
mysql> SELECT ST_AsText(ST_X(Point(56.7, 53.34), 10.5));
+-----------------------------------------+
| ST_AsText(ST_X(Point(56.7, 53.34), 10.5)) |
+-----------------------------------------+
| POINT(10.5 53.34)                       |
+-----------------------------------------+
```

- `ST_Y(p[, new_y_val])`

  With a single argument representing a valid `Point` object *p*, `ST_Y()` returns the Y-coordinate value of *p* as a double-precision number, or `NULL` if the argument is `NULL`. These conditions apply:

  - If the argument is a valid geometry but not a `Point` object, an `ER_UNEXPECTED_GEOMETRY_TYPE` error occurs.
  - If the argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  If the optional second argument representing a valid Y-coordinate value is given, `ST_Y()` returns a `Point` object with an Y coordinate equal to the second argument, or `NULL` if either argument is `NULL`. These conditions apply:

  - Error handling for the first argument is as for the single-argument syntax, as described previously.
  - If the coordinate value in the second argument is `-inf`, `+inf`, or `NaN`, an `ER_DATA_OUT_OF_RANGE` error occurs.

```
mysql> SELECT ST_Y(Point(56.7, 53.34));
+--------------------------+
| ST_Y(Point(56.7, 53.34)) |
+--------------------------+
|                    53.34 |
+--------------------------+
mysql> SELECT ST_AsText(ST_Y(Point(56.7, 53.34), 10.5));
+-----------------------------------------+
| ST_AsText(ST_Y(Point(56.7, 53.34), 10.5)) |
+-----------------------------------------+
| POINT(56.7 10.5)                        |
+-----------------------------------------+
```

12.15.7.3 LineString and MultiLineString Property Functions

A `LineString` consists of `Point` values. You can extract particular points of a `LineString`, count the number of points that it contains, or obtain its length.

Some functions in this section also work for `MultiLineString` values.

- `ST_EndPoint(ls)`

  Returns the `Point` that is the endpoint of the `LineString` value *ls*.

```
mysql> SET @ls = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_AsText(ST_EndPoint(ST_GeomFromText(@ls)));
+-------------------------------------------+
| ST_AsText(ST_EndPoint(ST_GeomFromText(@ls))) |
+-------------------------------------------+
| POINT(3 3)                                |
+-------------------------------------------+
```

- `ST_IsClosed(ls)`

  For a `LineString` value *`ls`*, `ST_IsClosed()` returns 1 if *`ls`* is closed (that is, its `ST_StartPoint()` and `ST_EndPoint()` values are the same).

  For a `MultiLineString` value *`ls`*, `ST_IsClosed()` returns 1 if *`ls`* is closed (that is, the `ST_StartPoint()` and `ST_EndPoint()` values are the same for each `LineString` in *`ls`*).

  `ST_IsClosed()` returns 0 if *`ls`* is not closed, and `NULL` if *`ls`* is `NULL`.

```
mysql> SET @ls1 = 'LineString(1 1,2 2,3 3,2 2)';
mysql> SET @ls2 = 'LineString(1 1,2 2,3 3,1 1)';

mysql> SELECT ST_IsClosed(ST_GeomFromText(@ls1));
+-----------------------------------+
| ST_IsClosed(ST_GeomFromText(@ls1)) |
+-----------------------------------+
|                                 0 |
+-----------------------------------+

mysql> SELECT ST_IsClosed(ST_GeomFromText(@ls2));
+-----------------------------------+
| ST_IsClosed(ST_GeomFromText(@ls2)) |
+-----------------------------------+
|                                 1 |
+-----------------------------------+

mysql> SET @ls3 = 'MultiLineString((1 1,2 2,3 3),(4 4,5 5))';

mysql> SELECT ST_IsClosed(ST_GeomFromText(@ls3));
+-----------------------------------+
| ST_IsClosed(ST_GeomFromText(@ls3)) |
+-----------------------------------+
|                                 0 |
+-----------------------------------+
```

- `ST_Length(ls)`

  Returns a double-precision number indicating the length of the `LineString` or `MultiLineString` value *`ls`* in its associated spatial reference. The length of a `MultiLineString` value is equal to the sum of the lengths of its elements.

```
mysql> SET @ls = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_Length(ST_GeomFromText(@ls));
+---------------------------------+
| ST_Length(ST_GeomFromText(@ls)) |
+---------------------------------+
|             2.8284271247461903 |
+---------------------------------+

mysql> SET @mls = 'MultiLineString((1 1,2 2,3 3),(4 4,5 5))';
mysql> SELECT ST_Length(ST_GeomFromText(@mls));
+----------------------------------+
| ST_Length(ST_GeomFromText(@mls)) |
+----------------------------------+
|                4.242640687119286 |
+----------------------------------+
```

- `ST_NumPoints(ls)`

  Returns the number of `Point` objects in the `LineString` value *`ls`*.

```
mysql> SET @ls = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_NumPoints(ST_GeomFromText(@ls));
+-----------------------------------+
| ST_NumPoints(ST_GeomFromText(@ls)) |
```

```
+-----------------------------------+
|                                 3 |
+-----------------------------------+
```

- `ST_PointN(ls, N)`

  Returns the *N*-th `Point` in the `Linestring` value *ls*. Points are numbered beginning with 1.

```
mysql> SET @ls = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_AsText(ST_PointN(ST_GeomFromText(@ls),2));
+--------------------------------------------+
| ST_AsText(ST_PointN(ST_GeomFromText(@ls),2)) |
+--------------------------------------------+
| POINT(2 2)                                 |
+--------------------------------------------+
```

- `ST_StartPoint(ls)`

  Returns the `Point` that is the start point of the `LineString` value *ls*.

```
mysql> SET @ls = 'LineString(1 1,2 2,3 3)';
mysql> SELECT ST_AsText(ST_StartPoint(ST_GeomFromText(@ls)));
+-----------------------------------------------+
| ST_AsText(ST_StartPoint(ST_GeomFromText(@ls))) |
+-----------------------------------------------+
| POINT(1 1)                                    |
+-----------------------------------------------+
```

12.15.7.4 Polygon and MultiPolygon Property Functions

These functions return properties of `Polygon` or `MultiPolygon` values.

- `ST_Area(poly)`

  Returns a double-precision number indicating the area of the argument, as measured in its spatial reference system. For arguments of dimension 0 or 1, the result is 0.

  The result is the sum of the area values of all components for a geometry collection. If a geometry collection is empty, its area is returned as 0.

```
mysql> SET @poly = 'Polygon((0 0,0 3,3 0,0 0),(1 1,1 2,2 1,1 1))';
mysql> SELECT ST_Area(ST_GeomFromText(@poly));
+---------------------------------+
| ST_Area(ST_GeomFromText(@poly)) |
+---------------------------------+
|                               4 |
+---------------------------------+

mysql> SET @mpoly =
    -> 'MultiPolygon(((0 0,0 3,3 3,3 0,0 0),(1 1,1 2,2 2,2 1,1 1)))';
mysql> SELECT ST_Area(ST_GeomFromText(@mpoly));
+----------------------------------+
| ST_Area(ST_GeomFromText(@mpoly)) |
+----------------------------------+
|                                8 |
+----------------------------------+
```

- `ST_Centroid(mpoly)`

  Returns the mathematical centroid for the `MultiPolygon` value *mpoly* as a `Point`. The result is not guaranteed to be on the `MultiPolygon`.

  This function processes geometry collections by computing the centroid point for components of highest dimension in the collection. Such components are extracted and made into a single

`MultiPolygon`, `MultiLineString`, or `MultiPoint` for centroid computation. If the argument is an empty geometry collection, the return value is `NULL`.

```
mysql> SET @poly =
    -> ST_GeomFromText('POLYGON((0 0,10 0,10 10,0 10,0 0),(5 5,7 5,7 7,5 7,5 5))');
mysql> SELECT ST_GeometryType(@poly),ST_AsText(ST_Centroid(@poly));
+------------------------+-------------------------------------------+
| ST_GeometryType(@poly) | ST_AsText(ST_Centroid(@poly))             |
+------------------------+-------------------------------------------+
| POLYGON                | POINT(4.958333333333333 4.958333333333333) |
+------------------------+-------------------------------------------+
```

- `ST_ExteriorRing(poly)`

  Returns the exterior ring of the `Polygon` value *`poly`* as a `LineString`.

```
mysql> SET @poly =
    -> 'Polygon((0 0,0 3,3 3,3 0,0 0),(1 1,1 2,2 2,2 1,1 1))';
mysql> SELECT ST_AsText(ST_ExteriorRing(ST_GeomFromText(@poly)));
+-------------------------------------------------+
| ST_AsText(ST_ExteriorRing(ST_GeomFromText(@poly))) |
+-------------------------------------------------+
| LINESTRING(0 0,0 3,3 3,3 0,0 0)                 |
+-------------------------------------------------+
```

- `ST_InteriorRingN(poly, N)`

  Returns the *`N`*-th interior ring for the `Polygon` value *`poly`* as a `LineString`. Rings are numbered beginning with 1.

```
mysql> SET @poly =
    -> 'Polygon((0 0,0 3,3 3,3 0,0 0),(1 1,1 2,2 2,2 1,1 1))';
mysql> SELECT ST_AsText(ST_InteriorRingN(ST_GeomFromText(@poly),1));
+-----------------------------------------------------+
| ST_AsText(ST_InteriorRingN(ST_GeomFromText(@poly),1)) |
+-----------------------------------------------------+
| LINESTRING(1 1,1 2,2 2,2 1,1 1)                     |
+-----------------------------------------------------+
```

- `ST_NumInteriorRing(poly)`, `ST_NumInteriorRings(poly)`

  Returns the number of interior rings in the `Polygon` value *`poly`*.

```
mysql> SET @poly =
    -> 'Polygon((0 0,0 3,3 3,3 0,0 0),(1 1,1 2,2 2,2 1,1 1))';
mysql> SELECT ST_NumInteriorRings(ST_GeomFromText(@poly));
+-------------------------------------------+
| ST_NumInteriorRings(ST_GeomFromText(@poly)) |
+-------------------------------------------+
|                                         1 |
+-------------------------------------------+
```

12.15.7.5 GeometryCollection Property Functions

These functions return properties of `GeometryCollection` values.

- `ST_GeometryN(gc, N)`

  Returns the *`N`*-th geometry in the `GeometryCollection` value *`gc`*. Geometries are numbered beginning with 1.

```
mysql> SET @gc = 'GeometryCollection(Point(1 1),LineString(2 2, 3 3))';
mysql> SELECT ST_AsText(ST_GeometryN(ST_GeomFromText(@gc),1));
+-----------------------------------------------+
```

```
| ST_AsText(ST_GeometryN(ST_GeomFromText(@gc),1)) |
+-------------------------------------------------+
| POINT(1 1)                                      |
+-------------------------------------------------+
```

- `ST_NumGeometries(gc)`

  Returns the number of geometries in the `GeometryCollection` value *`gc`*.

```
mysql> SET @gc = 'GeometryCollection(Point(1 1),LineString(2 2, 3 3))';
mysql> SELECT ST_NumGeometries(ST_GeomFromText(@gc));
+---------------------------------------+
| ST_NumGeometries(ST_GeomFromText(@gc)) |
+---------------------------------------+
|                                     2 |
+---------------------------------------+
```

12.15.8 Spatial Operator Functions

OpenGIS proposes a number of functions that can produce geometries. They are designed to implement spatial operators.

These functions support all argument type combinations except those that are inapplicable according to the Open Geospatial Consortium specification.

- `ST_Buffer(g, d[, strategy1[, strategy2[, strategy3]]])`

  Returns a geometry that represents all points whose distance from the geometry value *`g`* is less than or equal to a distance of *`d`*, or `NULL` if any argument is `NULL`. The SRID of the geometry argument must be 0 because `ST_Buffer()` supports only the Cartesian coordinate system. If any geometry argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.

  If the geometry argument is empty, `ST_Buffer()` returns an empty geometry.

  If the distance is 0, `ST_Buffer()` returns the geometry argument unchanged:

```
mysql> SET @pt = ST_GeomFromText('POINT(0 0)');
mysql> SELECT ST_AsText(ST_Buffer(@pt, 0));
+------------------------------+
| ST_AsText(ST_Buffer(@pt, 0)) |
+------------------------------+
| POINT(0 0)                   |
+------------------------------+
```

  `ST_Buffer()` supports negative distances for `Polygon` and `MultiPolygon` values, and for geometry collections containing `Polygon` or `MultiPolygon` values. The result may be an empty geometry. An `ER_WRONG_ARGUMENTS` error occurs for `ST_Buffer()` with a negative distance for `Point`, `MultiPoint`, `LineString`, and `MultiLineString` values, and for geometry collections not containing any `Polygon` or `MultiPolygon` values.

  `ST_Buffer()` permits up to three optional strategy arguments following the distance argument. Strategies influence buffer computation. These arguments are byte string values produced by the `ST_Buffer_Strategy()` function, to be used for point, join, and end strategies:

  - Point strategies apply to `Point` and `MultiPoint` geometries. If no point strategy is specified, the default is `ST_Buffer_Strategy('point_circle', 32)`.
  - Join strategies apply to `LineString`, `MultiLineString`, `Polygon`, and `MultiPolygon` geometries. If no join strategy is specified, the default is `ST_Buffer_Strategy('join_round', 32)`.
  - End strategies apply to `LineString` and `MultiLineString` geometries. If no end strategy is specified, the default is `ST_Buffer_Strategy('end_round', 32)`.

Up to one strategy of each type may be specified, and they may be given in any order. If multiple strategies of a given type are specified, an `ER_WRONG_ARGUMENTS` error occurs.

```
mysql> SET @pt = ST_GeomFromText('POINT(0 0)');
mysql> SET @pt_strategy = ST_Buffer_Strategy('point_square');
mysql> SELECT ST_AsText(ST_Buffer(@pt, 2, @pt_strategy));
+-----------------------------------------+
| ST_AsText(ST_Buffer(@pt, 2, @pt_strategy)) |
+-----------------------------------------+
| POLYGON((-2 -2,2 -2,2 2,-2 2,-2 -2))       |
+-----------------------------------------+
```

```
mysql> SET @ls = ST_GeomFromText('LINESTRING(0 0,0 5,5 5)');
mysql> SET @end_strategy = ST_Buffer_Strategy('end_flat');
mysql> SET @join_strategy = ST_Buffer_Strategy('join_round', 10);
mysql> SELECT ST_AsText(ST_Buffer(@ls, 5, @end_strategy, @join_strategy))
+-------------------------------------------------------------+
| ST_AsText(ST_Buffer(@ls, 5, @end_strategy, @join_strategy))   |
+-------------------------------------------------------------+
| POLYGON((5 5,5 10,0 10,-3.5355339059327373 8.535533905932738, |
| -5 5,-5 0,0 0,5 0,5 5))                                       |
+-------------------------------------------------------------+
```

- `ST_Buffer_Strategy(`*`strategy`*`[, `*`points_per_circle`*`])`

  This function returns a strategy byte string for use with `ST_Buffer()` to influence buffer computation. The result is `NULL` if any argument is `NULL`. If any argument is invalid, an `ER_WRONG_ARGUMENTS` error occurs.

  Information about strategies is available at Boost.org.

  The first argument must be a string indicating a strategy option:

  - For point strategies, permitted values are `'point_circle'` and `'point_square'`.
  - For join strategies, permitted values are `'join_round'` and `'join_miter'`.
  - For end strategies, permitted values are `'end_round'` and `'end_flat'`.

  If the first argument is `'point_circle'`, `'join_round'`, `'join_miter'`, or `'end_round'`, the *`points_per_circle`* argument must be given as a positive numeric value. The maximum *`points_per_circle`* value is the value of the `max_points_in_geometry` system variable. If the first argument is `'point_square'` or `'end_flat'`, the *`points_per_circle`* argument must not be given or an `ER_WRONG_ARGUMENTS` error occurs.

  For examples, see the description of `ST_Buffer()`.

- `ST_ConvexHull(`*`g`*`)`

  Returns a geometry that represents the convex hull of the geometry value *`g`*.

  This function computes a geometry's convex hull by first checking whether its vertex points are colinear. The function returns a linear hull if so, a polygon hull otherwise. This function processes geometry collections by extracting all vertex points of all components of the collection, creating a `MultiPoint` value from them, and computing its convex hull. If the argument is an empty geometry collection, the return value is `NULL`.

```
mysql> SET @g = 'MULTIPOINT(5 0,25 0,15 10,15 25)';
mysql> SELECT ST_AsText(ST_ConvexHull(ST_GeomFromText(@g)));
+---------------------------------------------+
| ST_AsText(ST_ConvexHull(ST_GeomFromText(@g))) |
+---------------------------------------------+
```

```
| POLYGON((5 0,25 0,15 25,5 0))                |
+----------------------------------------------+
```

- `ST_Difference(g1, g2)`

  Returns a geometry that represents the point set difference of the geometry values *g1* and *g2*.

```
mysql> SET @g1 = Point(1,1), @g2 = Point(2,2);
mysql> SELECT ST_AsText(ST_Difference(@g1, @g2));
+-----------------------------------+
| ST_AsText(ST_Difference(@g1, @g2)) |
+-----------------------------------+
| POINT(1 1)                        |
+-----------------------------------+
```

- `ST_Intersection(g1, g2)`

  Returns a geometry that represents the point set intersection of the geometry values *g1* and *g2*.

```
mysql> SET @g1 = ST_GeomFromText('LineString(1 1, 3 3)');
mysql> SET @g2 = ST_GeomFromText('LineString(1 3, 3 1)');
mysql> SELECT ST_AsText(ST_Intersection(@g1, @g2));
+------------------------------------+
| ST_AsText(ST_Intersection(@g1, @g2)) |
+------------------------------------+
| POINT(2 2)                         |
+------------------------------------+
```

- `ST_SymDifference(g1, g2)`

  Returns a geometry that represents the point set symmetric difference of the geometry values *g1* and *g2*, which is defined as:

```
g1 symdifference g2 := (g1 union g2) difference (g1 intersection g2)
```

  Or, in function call notation:

```
ST_SymDifference(g1, g2) = ST_Difference(ST_Union(g1, g2), ST_Intersection(g1, g2))
```

```
mysql> SET @g1 = Point(1,1), @g2 = Point(2,2);
mysql> SELECT ST_AsText(ST_SymDifference(@g1, @g2));
+-------------------------------------+
| ST_AsText(ST_SymDifference(@g1, @g2)) |
+-------------------------------------+
| MULTIPOINT((1 1),(2 2))             |
+-------------------------------------+
```

- `ST_Union(g1, g2)`

  Returns a geometry that represents the point set union of the geometry values *g1* and *g2*.

```
mysql> SET @g1 = ST_GeomFromText('LineString(1 1, 3 3)');
mysql> SET @g2 = ST_GeomFromText('LineString(1 3, 3 1)');
mysql> SELECT ST_AsText(ST_Union(@g1, @g2));
+-------------------------------------+
| ST_AsText(ST_Union(@g1, @g2))       |
+-------------------------------------+
| MULTILINESTRING((1 1,3 3),(1 3,3 1)) |
+-------------------------------------+
```

In addition, Section 12.15.7, “Geometry Property Functions”, discusses several functions that construct new geometries from existing ones. See that section for descriptions of these functions:

- `ST_Envelope(g)`
- `ST_StartPoint(ls)`
- `ST_EndPoint(ls)`
- `ST_PointN(ls, N)`
- `ST_ExteriorRing(poly)`
- `ST_InteriorRingN(poly, N)`
- `ST_GeometryN(gc, N)`

12.15.9 Functions That Test Spatial Relations Between Geometry Objects

The functions described in this section take two geometries as arguments and return a qualitative or quantitative relation between them.

MySQL implements two sets of functions using function names defined by the OpenGIS specification. One set tests the relationship between two geometry values using precise object shapes, the other set uses object minimum bounding rectangles (MBRs).

There is also a MySQL-specific set of MBR-based functions available to test the relationship between two geometry values.

12.15.9.1 Spatial Relation Functions That Use Object Shapes

The OpenGIS specification defines the following functions to test the relationship between two geometry values `g1` and `g2`, using precise object shapes. The return values 1 and 0 indicate true and false, respectively, except for `ST_Distance()`, which returns distance values.

In general, these functions support all argument type combinations except those that are inapplicable according to the Open Geospatial Consortium specification.

Exception: `ST_Distance()` detects arguments in a geographic (ellipsoidal) spatial reference system (SRS). For arguments in a geographic SRS, the only permitted argument types are `Point` and `Point`, or `Point` and `MultiPoint` (in any argument order).

- `ST_Contains(g1, g2)`

  Returns 1 or 0 to indicate whether *`g1`* completely contains *`g2`*. This tests the opposite relationship as `ST_Within()`.

  `ST_Contains()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `ST_Crosses(g1, g2)`

  Two geometries *spatially cross* if their spatial relation has the following properties:

- The two geometries intersect.
- Their intersection results in a geometry that has a dimension that is one less than the maximum dimension of the two given geometries.
- Their intersection is not equal to either of the two given geometries.

This function returns 1 or 0 to indicate whether *g1* spatially crosses *g2*. If g1 is of dimension 2 (Polygon or MultiPolygon), or *g2* is of dimension 1 (Point or MultiPoint), the return value is NULL.

ST_Crosses() handles its arguments as follows:

- If any argument is NULL or an empty geometry, the return value is NULL.
- If any argument is not a syntactically well-formed geometry, an ER_GIS_INVALID_DATA error occurs.
- If the arguments do not have the same SRID, an ER_GIS_DIFFERENT_SRIDS error occurs.
- If the arguments are in an undefined SRS, an ER_SRS_NOT_FOUND error occurs.
- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-NULL.

- ST_Disjoint(*g1*, *g2*)

  Returns 1 or 0 to indicate whether *g1* is spatially disjoint from (does not intersect) *g2*.

  ST_Disjoint() handles its arguments as follows:

  - If any argument is NULL or an empty geometry, the return value is NULL.
  - If any argument is not a syntactically well-formed geometry, an ER_GIS_INVALID_DATA error occurs.
  - If the arguments do not have the same SRID, an ER_GIS_DIFFERENT_SRIDS error occurs.
  - If the arguments are in an undefined SRS, an ER_SRS_NOT_FOUND error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-NULL.

- ST_Distance(*g1*, *g2*)

  Returns the distance between *g1* and *g2*, measured in the length unit of the spatial reference system (SRS).

  This function processes geometry collections by returning the shortest distance among all combinations of the components of the two geometry arguments.

  ST_Distance() detects arguments in a geographic (ellipsoidal) spatial reference system and returns the geodetic distance on the ellipsoid in meters. The only permitted geographic argument types are Point and Point, or Point and MultiPoint (in any argument order). If called with other geometry type argument combinations in a geographic SRS, an ER_NOT_IMPLEMENTED_FOR_GEOGRAPHIC_SRS error occurs. Calculations for projected SRSs and SRID 0 remain the same.

`ST_Distance()` handles its arguments as follows:

- If any argument is `NULL` or an empty geometry, the return value is `NULL`.
- If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
- If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
- If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
- If any argument is geometrically invalid, either the result is an undefined distance (that is, it can be any number), or an error occurs.
- If an intermediate or final result produces NaN or a negative number, an `ER_GIS_INVALID_DATA` error occurs.
- Otherwise, the return value is non-`NULL`.

For geographic SRS geometry arguments in an acceptable combination, an error occurs if any argument has a longitude or latitude that is out of range:

- Longitude values must be in the range (−180, 180], or an `ER_LONGITUDE_OUT_OF_RANGE` error occurs.
- Latitude values must be in the range [−90, 90], or an `ER_LATITUDE_OUT_OF_RANGE` error occurs.

```
mysql> SET @g1 = Point(1,1);
mysql> SET @g2 = Point(2,2);
mysql> SELECT ST_Distance(@g1, @g2);
+----------------------+
| ST_Distance(@g1, @g2) |
+----------------------+
|    1.4142135623730951 |
+----------------------+

mysql> SET @g1 = ST_GeomFromText('POINT(1 1)', 4326);;
mysql> SET @g2 = ST_GeomFromText('POINT(2 2)', 4326);;
mysql> SELECT ST_Distance(@g1, @g2);
+----------------------+
| ST_Distance(@g1, @g2) |
+----------------------+
|     156874.3859490455 |
+----------------------+
```

For the special case of distance calculations on a sphere, see the `ST_Distance_Sphere()` function.

- `ST_Equals(`*`g1, g2`*`)`

  Returns 1 or 0 to indicate whether *`g1`* is spatially equal to *`g2`*.

  `ST_Equals()` handles its arguments as follows:

  - If any argument is `NULL`, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.

- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-`NULL`.

```
mysql> SET @g1 = Point(1,1), @g2 = Point(2,2);
mysql> SELECT ST_Equals(@g1, @g1), ST_Equals(@g1, @g2);
+---------------------+---------------------+
| ST_Equals(@g1, @g1) | ST_Equals(@g1, @g2) |
+---------------------+---------------------+
|                   1 |                   0 |
+---------------------+---------------------+
```

- `ST_Intersects(`*`g1`*`, `*`g2`*`)`

  Returns 1 or 0 to indicate whether *`g1`* spatially intersects *`g2`*.

  `ST_Intersects()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `ST_Overlaps(`*`g1`*`, `*`g2`*`)`

  Two geometries *spatially overlap* if they intersect and their intersection results in a geometry of the same dimension but not equal to either of the given geometries.

  This function returns 1 or 0 to indicate whether *`g1`* spatially overlaps *`g2`*. If the dimensions of the two geometries are not equal, the return value is `NULL`.

  `ST_Overlaps()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `ST_Touches(`*`g1`*`, `*`g2`*`)`

  Two geometries *spatially touch* if their interiors do not intersect, but the boundary of one of the geometries intersects either the boundary or the interior of the other.

This function returns 1 or 0 to indicate whether *g1* spatially touches *g2*. If both geometries are of dimension 0 (`Point` or `MultiPoint`), the return value is `NULL`.

`ST_Touches()` handles its arguments as follows:

- If any argument is `NULL` or an empty geometry, the return value is `NULL`.
- If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
- If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
- If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-`NULL`.

- `ST_Within(g1, g2)`

  Returns 1 or 0 to indicate whether *g1* is spatially within *g2*. This tests the opposite relationship as `ST_Contains()`.

  `ST_Within()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

12.15.9.2 MySQL-Specific Spatial Relation Functions That Use Minimum Bounding Rectangles (MBRs)

MySQL provides several MySQL-specific functions that test relations between minimum bounding rectangles of two geometries `g1` and `g2`. The return values 1 and 0 indicate true and false, respectively.

The bounding box of a point is interpreted as a point that is both boundary and interior.

The bounding box of a straight horizontal or vertical line is interpreted as a line where the interior of the line is also boundary. The endpoints are boundary points.

If any of the parameters are geometry collections, the interior, boundary, and exterior of those parameters are those of the union of all elements in the collection.

- `MBRContains(g1, g2)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangle of *g1* contains the minimum bounding rectangle of *g2*. This tests the opposite relationship as `MBRWithin()`.

  `MBRContains()` handles its arguments as follows:

- If any argument is `NULL` or an empty geometry, the return value is `NULL`.
- If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
- If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
- If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-`NULL`.

```
mysql> SET @g1 = ST_GeomFromText('Polygon((0 0,0 3,3 3,3 0,0 0))');
mysql> SET @g2 = ST_GeomFromText('Point(1 1)');
mysql> SELECT MBRContains(@g1,@g2), MBRWithin(@g2,@g1);
+---------------------+-------------------+
| MBRContains(@g1,@g2) | MBRWithin(@g2,@g1) |
+---------------------+-------------------+
|                   1 |                 1 |
+---------------------+-------------------+
```

- `MBRCoveredBy(`*`g1, g2`*`)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangle of *`g1`* is covered by the minimum bounding rectangle of *`g2`*. This tests the opposite relationship as `MBRCovers()`.

  `MBRCoveredBy()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

```
mysql> SET @g1 = ST_GeomFromText('Polygon((0 0,0 3,3 3,3 0,0 0))');
mysql> SET @g2 = ST_GeomFromText('Point(1 1)');
mysql> SELECT MBRCovers(@g1,@g2), MBRCoveredby(@g1,@g2);
+-------------------+----------------------+
| MBRCovers(@g1,@g2) | MBRCoveredby(@g1,@g2) |
+-------------------+----------------------+
|                 1 |                    0 |
+-------------------+----------------------+
mysql> SELECT MBRCovers(@g2,@g1), MBRCoveredby(@g2,@g1);
+-------------------+----------------------+
| MBRCovers(@g2,@g1) | MBRCoveredby(@g2,@g1) |
+-------------------+----------------------+
|                 0 |                    1 |
+-------------------+----------------------+
```

- `MBRCovers(`*`g1, g2`*`)`

Returns 1 or 0 to indicate whether the minimum bounding rectangle of *g1* covers the minimum bounding rectangle of *g2*. This tests the opposite relationship as `MBRCoveredBy()`. See the description of `MBRCoveredBy()` for examples.

`MBRCovers()` handles its arguments as follows:

- If any argument is `NULL` or an empty geometry, the return value is `NULL`.
- If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
- If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
- If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-`NULL`.

- `MBRDisjoint(g1, g2)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangles of the two geometries *g1* and *g2* are disjoint (do not intersect).

  `MBRDisjoint()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `MBREquals(g1, g2)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangles of the two geometries *g1* and *g2* are the same.

  `MBREquals()` handles its arguments as follows:

  - If any argument is `NULL`, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `MBRIntersects(g1, g2)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangles of the two geometries *g1* and *g2* intersect.

  `MBRIntersects()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `MBROverlaps(g1, g2)`

  Two geometries *spatially overlap* if they intersect and their intersection results in a geometry of the same dimension but not equal to either of the given geometries.

  This function returns 1 or 0 to indicate whether the minimum bounding rectangles of the two geometries *g1* and *g2* overlap.

  `MBROverlaps()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

- `MBRTouches(g1, g2)`

  Two geometries *spatially touch* if their interiors do not intersect, but the boundary of one of the geometries intersects either the boundary or the interior of the other.

  This function returns 1 or 0 to indicate whether the minimum bounding rectangles of the two geometries *g1* and *g2* touch.

  `MBRTouches()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.

- If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
- If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
- Otherwise, the return value is non-`NULL`.

- `MBRWithin(g1, g2)`

  Returns 1 or 0 to indicate whether the minimum bounding rectangle of *`g1`* is within the minimum bounding rectangle of *`g2`*. This tests the opposite relationship as `MBRContains()`.

  `MBRWithin()` handles its arguments as follows:

  - If any argument is `NULL` or an empty geometry, the return value is `NULL`.
  - If any argument is not a syntactically well-formed geometry, an `ER_GIS_INVALID_DATA` error occurs.
  - If the arguments do not have the same SRID, an `ER_GIS_DIFFERENT_SRIDS` error occurs.
  - If the arguments are in an undefined SRS, an `ER_SRS_NOT_FOUND` error occurs.
  - If any argument is geometrically invalid, either the result is true or false (it is undefined which), or an error occurs.
  - Otherwise, the return value is non-`NULL`.

```
mysql> SET @g1 = ST_GeomFromText('Polygon((0 0,0 3,3 3,3 0,0 0))');
mysql> SET @g2 = ST_GeomFromText('Polygon((0 0,0 5,5 5,5 0,0 0))');
mysql> SELECT MBRWithin(@g1,@g2), MBRWithin(@g2,@g1);
+--------------------+--------------------+
| MBRWithin(@g1,@g2) | MBRWithin(@g2,@g1) |
+--------------------+--------------------+
|                  1 |                  0 |
+--------------------+--------------------+
```

12.15.10 Spatial Geohash Functions

Geohash is a system for encoding latitude and longitude coordinates of arbitrary precision into a text string. Geohash values are strings that contain only characters chosen from `"0123456789bcdefghjkmnpqrstuvwxyz"`.

The functions in this section enable manipulation of geohash values, which provides applications the capabilities of importing and exporting geohash data, and of indexing and searching geohash values.

- `ST_GeoHash(longitude, latitude, max_length)`, `ST_GeoHash(point, max_length)`

  Returns a geohash string in the connection character set and collation. The result is `NULL` if any argument is `NULL`. An error occurs if any argument is invalid.

  For the first syntax, the *`longitude`* must be a number in the range [−180, 180], and the *`latitude`* must be a number in the range [−90, 90]. For the second syntax, a `POINT` value is required, where the X and Y coordinates are in the valid ranges for longitude and latitude, respectively.

  The resulting string is no longer than *`max_length`* characters, which has an upper limit of 100. The string might be shorter than *`max_length`* characters because the algorithm that creates the geohash value continues until it has created a string that is either an exact representation of the location or *`max_length`* characters, whichever comes first.

```
mysql> SELECT ST_GeoHash(180,0,10), ST_GeoHash(-180,-90,15);
+----------------------+------------------------+
| ST_GeoHash(180,0,10) | ST_GeoHash(-180,-90,15) |
+----------------------+------------------------+
| xbpbpbpbpb           | 000000000000000         |
+----------------------+------------------------+
```

- `ST_LatFromGeoHash(geohash_str)`

  Returns the latitude from a geohash string value, as a `DOUBLE` value in the range [−90, 90]. The result is `NULL` if any argument is `NULL`. An error occurs if the argument is invalid.

  The `ST_LatFromGeoHash()` decoding function reads no more than 433 characters from the *`geohash_str`* argument. That represents the upper limit on information in the internal representation of coordinate values. Characters past the 433rd are ignored, even if they are otherwise illegal and produce an error.

```
mysql> SELECT ST_LatFromGeoHash(ST_GeoHash(45,-20,10));
+-----------------------------------------+
| ST_LatFromGeoHash(ST_GeoHash(45,-20,10)) |
+-----------------------------------------+
|                                     -20 |
+-----------------------------------------+
```

- `ST_LongFromGeoHash(geohash_str)`

  Returns the longitude from a geohash string value, as a `DOUBLE` value in the range [−180, 180]. The result is `NULL` if any argument is `NULL`. An error occurs if the argument is invalid.

  The remarks in the description of `ST_LatFromGeoHash()` regarding the maximum number of characters processed from the *`geohash_str`* argument also apply to `ST_LongFromGeoHash()`.

```
mysql> SELECT ST_LongFromGeoHash(ST_GeoHash(45,-20,10));
+------------------------------------------+
| ST_LongFromGeoHash(ST_GeoHash(45,-20,10)) |
+------------------------------------------+
|                                       45 |
+------------------------------------------+
```

- `ST_PointFromGeoHash(geohash_str, srid)`

  Returns a `POINT` value containing the decoded geohash value, given a geohash string value. The X and Y coordinates of the point are the longitude in the range [−180, 180] and the latitude in the range [−90, 90], respectively. The *`srid`* value is an unsigned 32-bit integer. The result is `NULL` if any argument is `NULL`. An error occurs if any argument is invalid.

  The remarks in the description of `ST_LatFromGeoHash()` regarding the maximum number of characters processed from the *`geohash_str`* argument also apply to `ST_PointFromGeoHash()`.

```
mysql> SET @gh = ST_GeoHash(45,-20,10);
mysql> SELECT ST_AsText(ST_PointFromGeoHash(@gh,0));
+---------------------------------------+
| ST_AsText(ST_PointFromGeoHash(@gh,0)) |
+---------------------------------------+
| POINT(45 -20)                         |
+---------------------------------------+
```

12.15.11 Spatial GeoJSON Functions

This section describes functions for converting between GeoJSON documents and spatial values. GeoJSON is an open standard for encoding geometric/geographical features. For more information, see http://geojson.org. The functions discussed here follow GeoJSON specification revision 1.0.

GeoJSON supports the same geometric/geographic data types that MySQL supports. Feature and FeatureCollection objects are not supported, except that geometry objects are extracted from them. CRS support is limited to values that identify an SRID.

MySQL also supports a native `JSON` data type and a set of SQL functions to enable operations on JSON values. For more information, see Section 11.6, “The JSON Data Type”, and Section 12.16, “JSON Functions”.

- `ST_AsGeoJSON(g [, max_dec_digits [, options]])`

  Generates a GeoJSON object from the geometry *`g`*. The object string has the connection character set and collation.

  `max_dec_digits`, if specified, limits the number of decimal digits for coordinates and causes rounding of output. If not specified, this argument defaults to its maximum value of $2^{32} - 1$. The minimum is 0.

  `options`, if specified, is a bitmask. The following table shows the permitted flag values. If the geometry argument has an SRID of 0, no CRS object is produced even for those flag values that request one.

Flag Value	Meaning
0	No options. This is the default if *`options`* is not specified.
1	Add a bounding box to the output.
2	Add a short-format CRS URN to the output. The default format is a short format (`EPSG:`*`srid`*).
4	Add a long-format CRS URN (`urn:ogc:def:crs:EPSG::`*`srid`*). This flag overrides flag 2. For example, option values of 5 and 7 mean the same (add a bounding box and a long-format CRS URN).

  If any argument is `NULL`, the return value is `NULL`. If any non-`NULL` argument is invalid, an error occurs.

```
mysql> SELECT ST_AsGeoJSON(ST_GeomFromText('POINT(11.11111 12.22222)'),2);
+-------------------------------------------------------------+
| ST_AsGeoJSON(ST_GeomFromText('POINT(11.11111 12.22222)'),2) |
+-------------------------------------------------------------+
| {"type": "Point", "coordinates": [11.11, 12.22]}            |
+-------------------------------------------------------------+
```

- `ST_GeomFromGeoJSON(str [, options [, srid]])`

  Parses a string *`str`* representing a GeoJSON object and returns a geometry.

  `options`, if given, describes how to handle GeoJSON documents that contain geometries with coordinate dimensions higher than 2. The following table shows the permitted *`options`* values.

Option Value	Meaning
1	Reject the document and produce an error. This is the default if *`options`* is not specified.
2, 3, 4	Accept the document and strip off the coordinates for higher coordinate dimensions.

  `options` values of 2, 3, and 4 currently produce the same effect. If geometries with coordinate dimensions higher than 2 are supported in the future, these values will produce different effects.

  The *`srid`* argument, if given, must be a 32-bit unsigned integer. If not given, the geometry return value has an SRID of 4326.

GeoJSON geometry, feature, and feature collection objects may have a `crs` property. The parsing function parses named CRS URNs in the `urn:ogc:def:crs:EPSG::srid` and `EPSG:srid` namespaces, but not CRSs given as link objects. Also, `urn:ogc:def:crs:OGC:1.3:CRS84` is recognized as SRID 4326. If an object has a CRS that is not understood, an error occurs, with the exception that if the optional *`srid`* argument is given, any CRS is ignored even if it is invalid.

As specified in the GeoJSON specification, parsing is case sensitive for the `type` member of the GeoJSON input (`Point`, `LineString`, and so forth). The specification is silent regarding case sensitivity for other parsing, which in MySQL is not case sensitive.

If any argument is `NULL`, the return value is `NULL`. If any non-`NULL` argument is invalid, an error occurs.

This example shows the parsing result for a simple GeoJSON object:

```
mysql> SET @json = '{ "type": "Point", "coordinates": [102.0, 0.0]}';
mysql> SELECT ST_AsText(ST_GeomFromGeoJSON(@json));
+---------------------------------------+
| ST_AsText(ST_GeomFromGeoJSON(@json)) |
+---------------------------------------+
| POINT(102 0)                          |
+---------------------------------------+
```

12.15.12 Spatial Convenience Functions

The functions in this section provide convenience operations on geometry values.

- `ST_Distance_Sphere(`*`g1, g2 [, radius`*`]`)`

  Returns the mimimum spherical distance between `Point` or `MultiPoint` arguments on a sphere, in meters. If any geometry argument is `NULL` or an empty geometry, the return value is `NULL`.

  Supported geometry argument combinations are `Point` and `Point`, or `Point` and `MultiPoint` (in any argument order). For other combinations, an `ER_GIS_UNSUPPORTED_ARGUMENT` error occurs.

  Calculations use a spherical earth and a configurable radius. The optional *`radius`* argument should be given in meters. If omitted, the default radius is 6,370,986 meters. If the *`radius`* argument is present but not positive, an `ER_WRONG_ARGUMENTS` error occurs.

  The geometry arguments should consist of points that specify (longitude, latitude) coordinate values:

  - Longitude and latitude are the first and second coordinates of the point, respectively.
  - Both coordinates are in degrees.
  - Longitude values must be in the range (−180, 180]. Positive values are east of the prime meridian.
  - Latitude values must be in the range [−90, 90]. Positive values are north of the equator.

  If any geometry argument is not a syntactically well-formed geometry byte string, an `ER_GIS_INVALID_DATA` error occurs.

```
mysql> SET @pt1 = ST_GeomFromText('POINT(0 0)');
mysql> SET @pt2 = ST_GeomFromText('POINT(180 0)');
mysql> SELECT ST_Distance_Sphere(@pt1, @pt2);
+--------------------------------+
| ST_Distance_Sphere(@pt1, @pt2) |
+--------------------------------+
|             20015042.813723423 |
+--------------------------------+
```

For general-purpose distance calculations, see the `ST_Distance()` function.

- `ST_IsValid(g)`

  Checks whether a geometry is valid, as defined by the OGC specification. `ST_IsValid()` returns 1 if the argument is a valid geometry byte string and is geometrically valid, 0 if the argument is not a valid geometry byte string or is not geometrically valid. If the argument is `NULL`, the return value is `NULL`.

  The only valid empty geometry is represented in the form of an empty geometry collection value. `ST_IsValid()` returns 1 in this case.

  `ST_IsValid()` works only for the Cartesian coordinate system and requires a geometry argument with an SRID of 0. An `ER_WRONG_ARGUMENTS` error occurs otherwise.

```
mysql> SET @ls1 = ST_GeomFromText('LINESTRING(0 0)');
mysql> SET @ls2 = ST_GeomFromText('LINESTRING(0 0, 1 1)');
mysql> SELECT ST_IsValid(@ls1);
+------------------+
| ST_IsValid(@ls1) |
+------------------+
|                0 |
+------------------+
mysql> SELECT ST_IsValid(@ls2);
+------------------+
| ST_IsValid(@ls2) |
+------------------+
|                1 |
+------------------+
```

- `ST_MakeEnvelope(pt1, pt2)`

  Returns the rectangle that forms the envelope around two points, as a `Point`, `LineString`, or `Polygon`. If any argument is `NULL`, the return value is `NULL`.

  Calculations are done using the Cartesian coordinate system rather than on a sphere, spheroid, or on earth.

  Given two points *`pt1`* and *`pt2`*, `ST_MakeEnvelope()` creates the result geometry on an abstract plane like this:

  - If *`pt1`* and *`pt2`* are equal, the result is the point *`pt1`*.
  - Otherwise, if *`(pt1, pt2)`* is a vertical or horizontal line segment, the result is the line segment *`(pt1, pt2)`*.
  - Otherwise, the result is a polygon using *`pt1`* and *`pt2`* as diagonal points.

  The result geometry has an SRID of 0.

  `ST_MakeEnvelope()` requires `Point` geometry arguments with an SRID of 0. An `ER_WRONG_ARGUMENTS` error occurs otherwise.

  If any argument is not a syntactically well-formed geometry byte string, or if any coordinate value of the two points is infinite (that is, NaN), an `ER_GIS_INVALID_DATA` error occurs.

```
mysql> SET @pt1 = ST_GeomFromText('POINT(0 0)');
mysql> SET @pt2 = ST_GeomFromText('POINT(1 1)');
mysql> SELECT ST_AsText(ST_MakeEnvelope(@pt1, @pt2));
+--------------------------------------+
| ST_AsText(ST_MakeEnvelope(@pt1, @pt2)) |
+--------------------------------------+
| POLYGON((0 0,1 0,1 1,0 1,0 0))         |
```

```
+------------------------------------+
```

- `ST_Simplify(g, max_distance)`

  Simplifies a geometry using the Douglas-Peucker algorithm and returns a simplified value of the same type. If any argument is `NULL`, the return value is `NULL`.

  The geometry may be any geometry type, although the Douglas-Peucker algorithm may not actually process every type. A geometry collection is processed by giving its components one by one to the simplification algorithm, and the returned geometries are put into a geometry collection as result.

  The *`max_distance`* argument is the distance (in units of the input coordinates) of a vertex to other segments to be removed. Vertices within this distance of the simplified linestring are removed. If the *`max_distance`* argument is not positive, or is NaN, an `ER_WRONG_ARGUMENTS` error occurs.

  According to Boost.Geometry, geometries might become invalid as a result of the simplification process, and the process might create self-intersections. To check the validity of the result, pass it to `ST_IsValid()`.

  If the geometry argument is not a syntactically well-formed geometry byte string, an `ER_GIS_INVALID_DATA` error occurs.

```
mysql> SET @g = ST_GeomFromText('LINESTRING(0 0,0 1,1 1,1 2,2 2,2 3,3 3)');
mysql> SELECT ST_AsText(ST_Simplify(@g, 0.5));
+---------------------------------+
| ST_AsText(ST_Simplify(@g, 0.5)) |
+---------------------------------+
| LINESTRING(0 0,0 1,1 1,2 3,3 3) |
+---------------------------------+
mysql> SELECT ST_AsText(ST_Simplify(@g, 1.0));
+---------------------------------+
| ST_AsText(ST_Simplify(@g, 1.0)) |
+---------------------------------+
| LINESTRING(0 0,3 3)             |
+---------------------------------+
mysql> SELECT ST_AsText(ST_Simplify(@g));
```

- `ST_Validate(g)`

  Validates a geometry according to the OGC specification. `ST_Validate()` returns the geometry if it is a valid geometry byte string and is geometrically valid, `NULL` if the argument is not a valid geometry byte string or is not geometrically valid or is `NULL`.

  A geometry can be a valid geometry byte string (WKB value plus SRID) but geometrically invalid. For example, this polygon is geometrically invalid: `POLYGON((0 0, 0 0, 0 0, 0 0, 0 0))`

  `ST_Validate()` can be used to filter out invalid geometry data, although at a cost. For applications that require more precise results not tainted by invalid data, this penalty may be worthwhile.

  If the geometry argument is valid, it is returned as is, except that if an input `Polygon` or `MultiPolygon` has clockwise rings, those rings are reversed before checking for validity. If the geometry is valid, the value with the reversed rings is returned.

  The only valid empty geometry is represented in the form of an empty geometry collection value. `ST_Validate()` returns it directly without further checks in this case.

  `ST_Validate()` works only for the Cartesian coordinate system and requires a geometry argument with an SRID of 0. An `ER_WRONG_ARGUMENTS` error occurs otherwise.

```
mysql> SET @ls1 = ST_GeomFromText('LINESTRING(0 0)');
mysql> SET @ls2 = ST_GeomFromText('LINESTRING(0 0, 1 1)');
mysql> SELECT ST_AsText(ST_Validate(@ls1));
+----------------------------+
```

```
| ST_AsText(ST_Validate(@ls1)) |
+------------------------------+
| NULL                         |
+------------------------------+
mysql> SELECT ST_AsText(ST_Validate(@ls2));
+------------------------------+
| ST_AsText(ST_Validate(@ls2)) |
+------------------------------+
| LINESTRING(0 0,1 1)          |
+------------------------------+
```

12.16 JSON Functions

The functions described in this section perform operations on JSON values. For discussion of the `JSON` data type and additional examples showing how to use these functions, see Section 11.6, "The JSON Data Type".

For functions that take a JSON argument, an error occurs if the argument is not a valid JSON value. Arguments parsed as JSON are indicated by *`json_doc`*; arguments indicated by *`val`* are not parsed.

A set of spatial functions for operating on GeoJSON values is also available. See Section 12.15.11, "Spatial GeoJSON Functions".

12.16.1 JSON Function Reference

Table 12.20 JSON Functions

Name	Description
`JSON_ARRAY()`	Create JSON array
`JSON_ARRAY_APPEND()`	Append data to JSON document
`JSON_ARRAY_INSERT()`	Insert into JSON array
`->`	Return value from JSON column after evaluating path; equivalent to JSON_EXTRACT().
`JSON_CONTAINS()`	Whether JSON document contains specific object at path
`JSON_CONTAINS_PATH()`	Whether JSON document contains any data at path
`JSON_DEPTH()`	Maximum depth of JSON document
`JSON_EXTRACT()`	Return data from JSON document
`->>`	Return value from JSON column after evaluating path and unquoting the result; equivalent to JSON_UNQUOTE(JSON_EXTRACT()).
`JSON_INSERT()`	Insert data into JSON document
`JSON_KEYS()`	Array of keys from JSON document
`JSON_LENGTH()`	Number of elements in JSON document
`JSON_MERGE()`	Merge JSON documents
`JSON_OBJECT()`	Create JSON object
`JSON_PRETTY()`	Prints a JSON document in human-readable format, with each array element or object member printed on a new line, indented two spaces with respect to its parent.
`JSON_QUOTE()`	Quote JSON document
`JSON_REMOVE()`	Remove data from JSON document
`JSON_REPLACE()`	Replace values in JSON document
`JSON_SEARCH()`	Path to value within JSON document
`JSON_SET()`	Insert data into JSON document

Name	Description
JSON_STORAGE_FREE()	Freed space within binary representation of a JSON column value following a partial update
JSON_STORAGE_SIZE()	Space used for storage of binary representation of a JSON document; for a JSON column, the space used when the document was inserted, prior to any partial updates
JSON_TYPE()	Type of JSON value
JSON_UNQUOTE()	Unquote JSON value
JSON_VALID()	Whether JSON value is valid

MySQL supports two aggregate JSON functions JSON_ARRAYAGG() and JSON_OBJECTAGG(). See Section 12.18, “Aggregate (GROUP BY) Functions”, for descriptions of these.

MySQL also supports “pretty-printing” of JSON values in an easy-to-read format, using the JSON_PRETTY() function. Beginning with MySQL 8.0.2, you can see how much storage space a given JSON value takes up, and how much space remains for additional storage, using JSON_STORAGE_SIZE() and JSON_STORAGE_FREE(), respectively. For complete descriptions of these functions, see Section 12.16.6, “JSON Utility Functions”.

12.16.2 Functions That Create JSON Values

The functions listed in this section compose JSON values from component elements.

- JSON_ARRAY([*val*[, *val*] ...])

  Evaluates a (possibly empty) list of values and returns a JSON array containing those values.

```
mysql> SELECT JSON_ARRAY(1, "abc", NULL, TRUE, CURTIME());
+---------------------------------------------+
| JSON_ARRAY(1, "abc", NULL, TRUE, CURTIME()) |
+---------------------------------------------+
| [1, "abc", null, true, "11:30:24.000000"]   |
+---------------------------------------------+
```

- JSON_OBJECT([*key*, *val*[, *key*, *val*] ...])

  Evaluates a (possibly empty) list of key-value pairs and returns a JSON object containing those pairs. An error occurs if any key name is NULL or the number of arguments is odd.

```
mysql> SELECT JSON_OBJECT('id', 87, 'name', 'carrot');
+-----------------------------------------+
| JSON_OBJECT('id', 87, 'name', 'carrot') |
+-----------------------------------------+
| {"id": 87, "name": "carrot"}            |
+-----------------------------------------+
```

- JSON_QUOTE(*string*)

  Quotes a string as a JSON value by wrapping it with double quote characters and escaping interior quote and other characters, then returning the result as a utf8mb4 string. Returns NULL if the argument is NULL.

  This function is typically used to produce a valid JSON string literal for inclusion within a JSON document.

  Certain special characters are escaped with backslashes per the escape sequences shown in Table 12.21, “JSON_UNQUOTE() Special Character Escape Sequences”.

```
mysql> SELECT JSON_QUOTE('null'), JSON_QUOTE('"null"');
```

```
+--------------------+----------------------+
| JSON_QUOTE('null') | JSON_QUOTE('"null"') |
+--------------------+----------------------+
| "null"             | "\"null\""           |
+--------------------+----------------------+
mysql> SELECT JSON_QUOTE('[1, 2, 3]');
+-------------------------+
| JSON_QUOTE('[1, 2, 3]') |
+-------------------------+
| "[1, 2, 3]"             |
+-------------------------+
```

You can also obtain JSON values by casting values of other types to the `JSON` type using `CAST(value AS JSON)`; see Converting between JSON and non-JSON values, for more information.

Two aggregate functions generating JSON values are available. `JSON_ARRAYAGG()` returns a result set as a single JSON array, and `JSON_OBJECTAGG()` returns a result set as a single JSON object. For more information, see Section 12.18, “Aggregate (GROUP BY) Functions”.

12.16.3 Functions That Search JSON Values

The functions in this section perform search operations on JSON values to extract data from them, report whether data exists at a location within them, or report the path to data within them.

- `JSON_CONTAINS(json_doc, val[, path])`

  Returns 0 or 1 to indicate whether a specific value is contained in a target JSON document, or, if a *`path`* argument is given, at a specific path within the target document. Returns `NULL` if any argument is `NULL` or the path argument does not identify a section of the target document. An error occurs if either document argument is not a valid JSON document or the *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  To check only whether any data exists at the path, use `JSON_CONTAINS_PATH()` instead.

  The following rules define containment:

  - A candidate scalar is contained in a target scalar if and only if they are comparable and are equal. Two scalar values are comparable if they have the same `JSON_TYPE()` types, with the exception that values of types `INTEGER` and `DECIMAL` are also comparable to each other.
  - A candidate array is contained in a target array if and only if every element in the candidate is contained in some element of the target.
  - A candidate nonarray is contained in a target array if and only if the candidate is contained in some element of the target.
  - A candidate object is contained in a target object if and only if for each key in the candidate there is a key with the same name in the target and the value associated with the candidate key is contained in the value associated with the target key.

  Otherwise, the candidate value is not contained in the target document.

```
mysql> SET @j = '{"a": 1, "b": 2, "c": {"d": 4}}';
mysql> SET @j2 = '1';
mysql> SELECT JSON_CONTAINS(@j, @j2, '$.a');
+-------------------------------+
| JSON_CONTAINS(@j, @j2, '$.a') |
+-------------------------------+
|                             1 |
+-------------------------------+
mysql> SELECT JSON_CONTAINS(@j, @j2, '$.b');
+-------------------------------+
| JSON_CONTAINS(@j, @j2, '$.b') |
+-------------------------------+
```

```
|                              0 |
+--------------------------------+

mysql> SET @j2 = '{"d": 4}';
mysql> SELECT JSON_CONTAINS(@j, @j2, '$.a');
+--------------------------------+
| JSON_CONTAINS(@j, @j2, '$.a') |
+--------------------------------+
|                              0 |
+--------------------------------+
mysql> SELECT JSON_CONTAINS(@j, @j2, '$.c');
+--------------------------------+
| JSON_CONTAINS(@j, @j2, '$.c') |
+--------------------------------+
|                              1 |
+--------------------------------+
```

- `JSON_CONTAINS_PATH(json_doc, one_or_all, path[, path] ...)`

  Returns 0 or 1 to indicate whether a JSON document contains data at a given path or paths. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document, any *`path`* argument is not a valid path expression, or *`one_or_all`* is not `'one'` or `'all'`.

  To check for a specific value at a path, use `JSON_CONTAINS()` instead.

  The return value is 0 if no specified path exists within the document. Otherwise, the return value depends on the *`one_or_all`* argument:

  - `'one'`: 1 if at least one path exists within the document, 0 otherwise.
  - `'all'`: 1 if all paths exist within the document, 0 otherwise.

```
mysql> SET @j = '{"a": 1, "b": 2, "c": {"d": 4}}';
mysql> SELECT JSON_CONTAINS_PATH(@j, 'one', '$.a', '$.e');
+-------------------------------------------+
| JSON_CONTAINS_PATH(@j, 'one', '$.a', '$.e') |
+-------------------------------------------+
|                                         1 |
+-------------------------------------------+
mysql> SELECT JSON_CONTAINS_PATH(@j, 'all', '$.a', '$.e');
+-------------------------------------------+
| JSON_CONTAINS_PATH(@j, 'all', '$.a', '$.e') |
+-------------------------------------------+
|                                         0 |
+-------------------------------------------+
mysql> SELECT JSON_CONTAINS_PATH(@j, 'one', '$.c.d');
+-------------------------------------+
| JSON_CONTAINS_PATH(@j, 'one', '$.c.d') |
+-------------------------------------+
|                                   1 |
+-------------------------------------+
mysql> SELECT JSON_CONTAINS_PATH(@j, 'one', '$.a.d');
+-------------------------------------+
| JSON_CONTAINS_PATH(@j, 'one', '$.a.d') |
+-------------------------------------+
|                                   0 |
+-------------------------------------+
```

- `JSON_EXTRACT(json_doc, path[, path] ...)`

  Returns data from a JSON document, selected from the parts of the document matched by the *`path`* arguments. Returns `NULL` if any argument is `NULL` or no paths locate a value in the document. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression.

The return value consists of all values matched by the *`path`* arguments. If it is possible that those arguments could return multiple values, the matched values are autowrapped as an array, in the order corresponding to the paths that produced them. Otherwise, the return value is the single matched value.

```
mysql> SELECT JSON_EXTRACT('[10, 20, [30, 40]]', '$[1]');
+-----------------------------------------+
| JSON_EXTRACT('[10, 20, [30, 40]]', '$[1]') |
+-----------------------------------------+
| 20                                      |
+-----------------------------------------+
mysql> SELECT JSON_EXTRACT('[10, 20, [30, 40]]', '$[1]', '$[0]');
+--------------------------------------------------+
| JSON_EXTRACT('[10, 20, [30, 40]]', '$[1]', '$[0]') |
+--------------------------------------------------+
| [20, 10]                                         |
+--------------------------------------------------+
mysql> SELECT JSON_EXTRACT('[10, 20, [30, 40]]', '$[2][*]');
+-----------------------------------------+
| JSON_EXTRACT('[10, 20, [30, 40]]', '$[2][*]') |
+-----------------------------------------+
| [30, 40]                                |
+-----------------------------------------+
```

MySQL supports the `->` operator as shorthand for this function as used with 2 arguments where the left hand side is a `JSON` column identifier (not an expression) and the right hand side is the JSON path to be matched within the column.

- *`column->path`*

  The `->` operator serves as an alias for the `JSON_EXTRACT()` function when used with two arguments, a column identifier on the left and a JSON path on the right that is evaluated against the JSON document (the column value). You can use such expressions in place of column identifiers wherever they occur in SQL statements.

  The two `SELECT` statements shown here produce the same output:

```
mysql> SELECT c, JSON_EXTRACT(c, "$.id"), g
    > FROM jemp
    > WHERE JSON_EXTRACT(c, "$.id") > 1
    > ORDER BY JSON_EXTRACT(c, "$.name");
+------------------------------+-----------+------+
| c                            | c->"$.id" | g    |
+------------------------------+-----------+------+
| {"id": "3", "name": "Barney"} | "3"       |    3 |
| {"id": "4", "name": "Betty"}  | "4"       |    4 |
| {"id": "2", "name": "Wilma"}  | "2"       |    2 |
+------------------------------+-----------+------+
3 rows in set (0.00 sec)

mysql> SELECT c, c->"$.id", g
    > FROM jemp
    > WHERE c->"$.id" > 1
    > ORDER BY c->"$.name";
+------------------------------+-----------+------+
| c                            | c->"$.id" | g    |
+------------------------------+-----------+------+
| {"id": "3", "name": "Barney"} | "3"       |    3 |
| {"id": "4", "name": "Betty"}  | "4"       |    4 |
| {"id": "2", "name": "Wilma"}  | "2"       |    2 |
+------------------------------+-----------+------+
3 rows in set (0.00 sec)
```

This functionality is not limited to `SELECT`, as shown here:

```
mysql> ALTER TABLE jemp ADD COLUMN n INT;
Query OK, 0 rows affected (0.68 sec)
Records: 0  Duplicates: 0  Warnings: 0

mysql> UPDATE jemp SET n=1 WHERE c->"$.id" = "4";
Query OK, 1 row affected (0.04 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT c, c->"$.id", g, n
     > FROM jemp
     > WHERE JSON_EXTRACT(c, "$.id") > 1
     > ORDER BY c->"$.name";
+------------------------------+-----------+------+------+
| c                            | c->"$.id" | g    | n    |
+------------------------------+-----------+------+------+
| {"id": "3", "name": "Barney"} | "3"       |    3 | NULL |
| {"id": "4", "name": "Betty"}  | "4"       |    4 |    1 |
| {"id": "2", "name": "Wilma"}  | "2"       |    2 | NULL |
+------------------------------+-----------+------+------+
3 rows in set (0.00 sec)

mysql> DELETE FROM jemp WHERE c->"$.id" = "4";
Query OK, 1 row affected (0.04 sec)

mysql> SELECT c, c->"$.id", g, n
     > FROM jemp
     > WHERE JSON_EXTRACT(c, "$.id") > 1
     > ORDER BY c->"$.name";
+------------------------------+-----------+------+------+
| c                            | c->"$.id" | g    | n    |
+------------------------------+-----------+------+------+
| {"id": "3", "name": "Barney"} | "3"       |    3 | NULL |
| {"id": "2", "name": "Wilma"}  | "2"       |    2 | NULL |
+------------------------------+-----------+------+------+
2 rows in set (0.00 sec)
```

(See Indexing a Generated Column to Provide a JSON Column Index, for the statements used to create and populate the table just shown.)

This also works with JSON array values, as shown here:

```
mysql> CREATE TABLE tj10 (a JSON, b INT);
Query OK, 0 rows affected (0.26 sec)

mysql> INSERT INTO tj10
     > VALUES ("[3,10,5,17,44]", 33), ("[3,10,5,17,[22,44,66]]", 0);
Query OK, 1 row affected (0.04 sec)

mysql> SELECT a->"$[4]" FROM tj10;
+--------------+
| a->"$[4]"    |
+--------------+
| 44           |
| [22, 44, 66] |
+--------------+
2 rows in set (0.00 sec)

mysql> SELECT * FROM tj10 WHERE a->"$[0]" = 3;
+------------------------------+------+
| a                            | b    |
+------------------------------+------+
| [3, 10, 5, 17, 44]           |   33 |
| [3, 10, 5, 17, [22, 44, 66]] |    0 |
+------------------------------+------+
2 rows in set (0.00 sec)
```

Nested arrays are supported. An expression using `->` evaluates as `NULL` if no matching key is found in the target JSON document, as shown here:

```
mysql> SELECT * FROM tj10 WHERE a->"$[4][1]" IS NOT NULL;
+-----------------------------+------+
| a                           | b    |
+-----------------------------+------+
| [3, 10, 5, 17, [22, 44, 66]] |    0 |
+-----------------------------+------+

mysql> SELECT a->"$[4][1]" FROM tj10;
+--------------+
| a->"$[4][1]" |
+--------------+
| NULL         |
| 44           |
+--------------+
2 rows in set (0.00 sec)
```

This is the same behavior as seen in such cases when using `JSON_EXTRACT()`:

```
mysql> SELECT JSON_EXTRACT(a, "$[4][1]") FROM tj10;
+----------------------------+
| JSON_EXTRACT(a, "$[4][1]") |
+----------------------------+
| NULL                       |
| 44                         |
+----------------------------+
2 rows in set (0.00 sec)
```

- *`column->>path`*

  This is an improved, unquoting extraction operator. Whereas the `->` operator simply extracts a value, the `->>` operator in addition unquotes the extracted result. In other words, given a `JSON` column value *`column`* and a path expression *`path`*, the following three expressions return the same value:

  - `JSON_UNQUOTE( JSON_EXTRACT(`*`column`*`, `*`path`*`) )`
  - `JSON_UNQUOTE(`*`column`*` -> `*`path`*`)`
  - *`column->>path`*

  The `->>` operator can be used wherever `JSON_UNQUOTE(JSON_EXTRACT())` would be allowed. This includes (but is not limited to) `SELECT` lists, `WHERE` and `HAVING` clauses, and `ORDER BY` and `GROUP BY` clauses.

  The next few statements demonstrate some `->>` operator equivalences with other expressions in the `mysql` client:

```
mysql> SELECT * FROM jemp WHERE g > 2;
+------------------------------+------+
| c                            | g    |
+------------------------------+------+
| {"id": "3", "name": "Barney"} |    3 |
| {"id": "4", "name": "Betty"}  |    4 |
+------------------------------+------+
2 rows in set (0.01 sec)

mysql> SELECT c->'$.name' AS name
    ->     FROM jemp WHERE g > 2;
+----------+
| name     |
+----------+
| "Barney" |
| "Betty"  |
+----------+
2 rows in set (0.00 sec)

mysql> SELECT JSON_UNQUOTE(c->'$.name') AS name
    ->     FROM jemp WHERE g > 2;
```

```
+--------+
| name   |
+--------+
| Barney |
| Betty  |
+--------+
2 rows in set (0.00 sec)

mysql> SELECT c->>'$.name' AS name
    ->     FROM jemp WHERE g > 2;
+--------+
| name   |
+--------+
| Barney |
| Betty  |
+--------+
2 rows in set (0.00 sec)
```

See Indexing a Generated Column to Provide a JSON Column Index, for the SQL statements used to create and populate the `jemp` table in the set of examples just shown.

This operator can also be used with JSON arrays, as shown here:

```
mysql> CREATE TABLE tj10 (a JSON, b INT);
Query OK, 0 rows affected (0.26 sec)

mysql> INSERT INTO tj10 VALUES
    ->     ('[3,10,5,"x",44]', 33),
    ->     ('[3,10,5,17,[22,"y",66]]', 0);
Query OK, 2 rows affected (0.04 sec)
Records: 2  Duplicates: 0  Warnings: 0

mysql> SELECT a->"$[3]", a->"$[4][1]" FROM tj10;
+-----------+--------------+
| a->"$[3]" | a->"$[4][1]" |
+-----------+--------------+
| "x"       | NULL         |
| 17        | "y"          |
+-----------+--------------+
2 rows in set (0.00 sec)

mysql> SELECT a->>"$[3]", a->>"$[4][1]" FROM tj10;
+------------+---------------+
| a->>"$[3]" | a->>"$[4][1]" |
+------------+---------------+
| x          | NULL          |
| 17         | y             |
+------------+---------------+
2 rows in set (0.00 sec)
```

As with `->`, the `->>` operator is always expanded in the output of `EXPLAIN`, as the following example demonstrates:

```
mysql> EXPLAIN SELECT c->>'$.name' AS name
    ->     FROM jemp WHERE g > 2\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: jemp
   partitions: NULL
         type: range
possible_keys: i
          key: i
      key_len: 5
          ref: NULL
         rows: 2
     filtered: 100.00
        Extra: Using where
1 row in set, 1 warning (0.00 sec)
```

```
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: /* select#1 */ select
json_unquote(json_extract(`jtest`.`jemp`.`c`,'$.name')) AS `name` from
`jtest`.`jemp` where (`jtest`.`jemp`.`g` > 2)
1 row in set (0.00 sec)
```

This is similar to how MySQL expands the `->` operator in the same circumstances.

- `JSON_KEYS(json_doc[, path])`

  Returns the keys from the top-level value of a JSON object as a JSON array, or, if a *`path`* argument is given, the top-level keys from the selected path. Returns `NULL` if any argument is `NULL`, the *`json_doc`* argument is not an object, or *`path`*, if given, does not locate an object. An error occurs if the *`json_doc`* argument is not a valid JSON document or the *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The result array is empty if the selected object is empty. If the top-level value has nested subobjects, the return value does not include keys from those subobjects.

```
mysql> SELECT JSON_KEYS('{"a": 1, "b": {"c": 30}}');
+---------------------------------------+
| JSON_KEYS('{"a": 1, "b": {"c": 30}}') |
+---------------------------------------+
| ["a", "b"]                            |
+---------------------------------------+
mysql> SELECT JSON_KEYS('{"a": 1, "b": {"c": 30}}', '$.b');
+----------------------------------------------+
| JSON_KEYS('{"a": 1, "b": {"c": 30}}', '$.b') |
+----------------------------------------------+
| ["c"]                                        |
+----------------------------------------------+
```

- `JSON_SEARCH(json_doc, one_or_all, search_str[, escape_char[, path] ...])`

  Returns the path to the given string within a JSON document. Returns `NULL` if any of the *`json_doc`*, *`search_str`*, or *`path`* arguments are `NULL`; no *`path`* exists within the document; or *`search_str`* is not found. An error occurs if the *`json_doc`* argument is not a valid JSON document, any *`path`* argument is not a valid path expression, *`one_or_all`* is not `'one'` or `'all'`, or *`escape_char`* is not a constant expression.

  The *`one_or_all`* argument affects the search as follows:

  - `'one'`: The search terminates after the first match and returns one path string. It is undefined which match is considered first.
  - `'all'`: The search returns all matching path strings such that no duplicate paths are included. If there are multiple strings, they are autowrapped as an array. The order of the array elements is undefined.

  Within the *`search_str`* search string argument, the `%` and `_` characters work as for the `LIKE` operator: `%` matches any number of characters (including zero characters), and `_` matches exactly one character.

  To specify a literal `%` or `_` character in the search string, precede it by the escape character. The default is `\` if the *`escape_char`* argument is missing or `NULL`. Otherwise, *`escape_char`* must be a constant that is empty or one character.

  For more information about matching and escape character behavior, see the description of `LIKE` in Section 12.5.1, “String Comparison Functions”. For escape character handling, a difference from

the `LIKE` behavior is that the escape character for `JSON_SEARCH()` must evaluate to a constant at compile time, not just at execution time. For example, if `JSON_SEARCH()` is used in a prepared statement and the *`escape_char`* argument is supplied using a `?` parameter, the parameter value might be constant at execution time, but is not at compile time.

```
mysql> SET @j = '["abc", [{"k": "10"}, "def"], {"x":"abc"}, {"y":"bcd"}]';

mysql> SELECT JSON_SEARCH(@j, 'one', 'abc');
+-------------------------------+
| JSON_SEARCH(@j, 'one', 'abc') |
+-------------------------------+
| "$[0]"                        |
+-------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', 'abc');
+-------------------------------+
| JSON_SEARCH(@j, 'all', 'abc') |
+-------------------------------+
| ["$[0]", "$[2].x"]            |
+-------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', 'ghi');
+-------------------------------+
| JSON_SEARCH(@j, 'all', 'ghi') |
+-------------------------------+
| NULL                          |
+-------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10');
+-------------------------------+
| JSON_SEARCH(@j, 'all', '10') |
+-------------------------------+
| "$[1][0].k"                   |
+-------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$');
+----------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$') |
+----------------------------------------+
| "$[1][0].k"                            |
+----------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$[*]');
+-----------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$[*]') |
+-----------------------------------------+
| "$[1][0].k"                             |
+-----------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$**.k');
+-----------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$**.k') |
+-----------------------------------------+
| "$[1][0].k"                             |
+-----------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$[*][0].k');
+-----------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$[*][0].k') |
+-----------------------------------------+
| "$[1][0].k"                             |
+-----------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$[1]');
+-----------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$[1]') |
+-----------------------------------------+
| "$[1][0].k"                             |
+-----------------------------------------+
```

```
mysql> SELECT JSON_SEARCH(@j, 'all', '10', NULL, '$[1][0]');
+-----------------------------------------------+
| JSON_SEARCH(@j, 'all', '10', NULL, '$[1][0]') |
+-----------------------------------------------+
| "$[1][0].k"                                   |
+-----------------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', 'abc', NULL, '$[2]');
+-------------------------------------------+
| JSON_SEARCH(@j, 'all', 'abc', NULL, '$[2]') |
+-------------------------------------------+
| "$[2].x"                                  |
+-------------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%a%');
+------------------------------+
| JSON_SEARCH(@j, 'all', '%a%') |
+------------------------------+
| ["$[0]", "$[2].x"]           |
+------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%');
+------------------------------+
| JSON_SEARCH(@j, 'all', '%b%') |
+------------------------------+
| ["$[0]", "$[2].x", "$[3].y"]  |
+------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%', NULL, '$[0]');
+-------------------------------------------+
| JSON_SEARCH(@j, 'all', '%b%', NULL, '$[0]') |
+-------------------------------------------+
| "$[0]"                                    |
+-------------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%', NULL, '$[2]');
+-------------------------------------------+
| JSON_SEARCH(@j, 'all', '%b%', NULL, '$[2]') |
+-------------------------------------------+
| "$[2].x"                                  |
+-------------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%', NULL, '$[1]');
+-------------------------------------------+
| JSON_SEARCH(@j, 'all', '%b%', NULL, '$[1]') |
+-------------------------------------------+
| NULL                                      |
+-------------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%', '', '$[1]');
+----------------------------------------+
| JSON_SEARCH(@j, 'all', '%b%', '', '$[1]') |
+----------------------------------------+
| NULL                                   |
+----------------------------------------+

mysql> SELECT JSON_SEARCH(@j, 'all', '%b%', '', '$[3]');
+----------------------------------------+
| JSON_SEARCH(@j, 'all', '%b%', '', '$[3]') |
+----------------------------------------+
| "$[3].y"                               |
+----------------------------------------+
```

For more information about the JSON path syntax supported by MySQL, including rules governing the wildcard operators `*` and `**`, see Section 12.16.7, “JSON Path Syntax”.

12.16.4 Functions That Modify JSON Values

The functions in this section modify JSON values and return the result.

- `JSON_APPEND(json_doc, path, val[, path, val] ...)`

This function was renamed to `JSON_ARRAY_APPEND()`, so use it instead.

- `JSON_ARRAY_APPEND(json_doc, path, val[, path, val] ...)`

  Appends values to the end of the indicated arrays within a JSON document and returns the result. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

  If a path selects a scalar or object value, that value is autowrapped within an array and the new value is added to that array. Pairs for which the path does not identify any value in the JSON document are ignored.

```
mysql> SET @j = '["a", ["b", "c"], "d"]';
mysql> SELECT JSON_ARRAY_APPEND(@j, '$[1]', 1);
+---------------------------------+
| JSON_ARRAY_APPEND(@j, '$[1]', 1) |
+---------------------------------+
| ["a", ["b", "c", 1], "d"]       |
+---------------------------------+
mysql> SELECT JSON_ARRAY_APPEND(@j, '$[0]', 2);
+---------------------------------+
| JSON_ARRAY_APPEND(@j, '$[0]', 2) |
+---------------------------------+
| [["a", 2], ["b", "c"], "d"]     |
+---------------------------------+
mysql> SELECT JSON_ARRAY_APPEND(@j, '$[1][0]', 3);
+------------------------------------+
| JSON_ARRAY_APPEND(@j, '$[1][0]', 3) |
+------------------------------------+
| ["a", [["b", 3], "c"], "d"]        |
+------------------------------------+

mysql> SET @j = '{"a": 1, "b": [2, 3], "c": 4}';
mysql> SELECT JSON_ARRAY_APPEND(@j, '$.b', 'x');
+-----------------------------------+
| JSON_ARRAY_APPEND(@j, '$.b', 'x')  |
+-----------------------------------+
| {"a": 1, "b": [2, 3, "x"], "c": 4} |
+-----------------------------------+
mysql> SELECT JSON_ARRAY_APPEND(@j, '$.c', 'y');
+-----------------------------------+
| JSON_ARRAY_APPEND(@j, '$.c', 'y')    |
+-----------------------------------+
| {"a": 1, "b": [2, 3], "c": [4, "y"]} |
+-----------------------------------+

mysql> SET @j = '{"a": 1}';
mysql> SELECT JSON_ARRAY_APPEND(@j, '$', 'z');
+---------------------------------+
| JSON_ARRAY_APPEND(@j, '$', 'z') |
+---------------------------------+
| [{"a": 1}, "z"]                 |
+---------------------------------+
```

- `JSON_ARRAY_INSERT(json_doc, path, val[, path, val] ...)`

  Updates a JSON document, inserting into an array within the document and returning the modified document. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard or does not end with an array element identifier.

  The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

Pairs for which the path does not identify any array in the JSON document are ignored. If a path identifies an array element, the corresponding value is inserted at that element position, shifting any following values to the right. If a path identifies an array position past the end of an array, the value is inserted at the end of the array.

```
mysql> SET @j = '["a", {"b": [1, 2]}, [3, 4]]';
mysql> SELECT JSON_ARRAY_INSERT(@j, '$[1]', 'x');
+-------------------------------------+
| JSON_ARRAY_INSERT(@j, '$[1]', 'x') |
+-------------------------------------+
| ["a", "x", {"b": [1, 2]}, [3, 4]]  |
+-------------------------------------+
mysql> SELECT JSON_ARRAY_INSERT(@j, '$[100]', 'x');
+-------------------------------------+
| JSON_ARRAY_INSERT(@j, '$[100]', 'x') |
+-------------------------------------+
| ["a", {"b": [1, 2]}, [3, 4], "x"]    |
+-------------------------------------+
mysql> SELECT JSON_ARRAY_INSERT(@j, '$[1].b[0]', 'x');
+---------------------------------------+
| JSON_ARRAY_INSERT(@j, '$[1].b[0]', 'x') |
+---------------------------------------+
| ["a", {"b": ["x", 1, 2]}, [3, 4]]       |
+---------------------------------------+
mysql> SELECT JSON_ARRAY_INSERT(@j, '$[2][1]', 'y');
+-------------------------------------+
| JSON_ARRAY_INSERT(@j, '$[2][1]', 'y') |
+-------------------------------------+
| ["a", {"b": [1, 2]}, [3, "y", 4]]     |
+-------------------------------------+
mysql> SELECT JSON_ARRAY_INSERT(@j, '$[0]', 'x', '$[2][1]', 'y');
+---------------------------------------------------+
| JSON_ARRAY_INSERT(@j, '$[0]', 'x', '$[2][1]', 'y') |
+---------------------------------------------------+
| ["x", "a", {"b": [1, 2]}, [3, 4]]                  |
+---------------------------------------------------+
```

Earlier modifications affect the positions of the following elements in the array, so subsequent paths in the same `JSON_ARRAY_INSERT()` call should take this into account. In the final example, the second path inserts nothing because the path no longer matches anything after the first insert.

- `JSON_INSERT(`*`json_doc`*`, `*`path`*`, `*`val`*`[, `*`path`*`, `*`val`*`] ...)`

  Inserts data into a JSON document and returns the result. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

  A path-value pair for an existing path in the document is ignored and does not overwrite the existing document value. A path-value pair for a nonexisting path in the document adds the value to the document if the path identifies one of these types of values:

  - A member not present in an existing object. The member is added to the object and associated with the new value.
  - A position past the end of an existing array. The array is extended with the new value. If the existing value is not an array, it is autowrapped as an array, then extended with the new value.

  Otherwise, a path-value pair for a nonexisting path in the document is ignored and has no effect.

  For a comparison of `JSON_INSERT()`, `JSON_REPLACE()`, and `JSON_SET()`, see the discussion of `JSON_SET()`.

```
mysql> SET @j = '{ "a": 1, "b": [2, 3]}';
mysql> SELECT JSON_INSERT(@j, '$.a', 10, '$.c', '[true, false]');
+-------------------------------------------------+
| JSON_INSERT(@j, '$.a', 10, '$.c', '[true, false]') |
+-------------------------------------------------+
| {"a": 1, "b": [2, 3], "c": "[true, false]"}        |
+-------------------------------------------------+
```

The third and final value listed in the result is a quoted string and not an array like the second one (which is not quoted in the output); no casting of values to the JSON type is performed. To insert the array as an array, you must perform such casts explicitly, as shown here:

```
mysql> SELECT JSON_INSERT(@j, '$.a', 10, '$.c', CAST('[true, false]' AS JSON));
+-----------------------------------------------------------------+
| JSON_INSERT(@j, '$.a', 10, '$.c', CAST('[true, false]' AS JSON)) |
+-----------------------------------------------------------------+
| {"a": 1, "b": [2, 3], "c": [true, false]}                       |
+-----------------------------------------------------------------+
1 row in set (0.00 sec)
```

- `JSON_MERGE(json_doc, json_doc[, json_doc] ...)`

  Merges two or more JSON documents and returns the merged result. Returns `NULL` if any argument is `NULL`. An error occurs if any argument is not a valid JSON document.

  Merging takes place according to the following rules. For additional information, see Normalization, Merging, and Autowrapping of JSON Values.

  - Adjacent arrays are merged to a single array.
  - Adjacent objects are merged to a single object.
  - A scalar value is autowrapped as an array and merged as an array.
  - An adjacent array and object are merged by autowrapping the object as an array and merging the two arrays.

```
mysql> SELECT JSON_MERGE('[1, 2]', '[true, false]');
+---------------------------------------+
| JSON_MERGE('[1, 2]', '[true, false]') |
+---------------------------------------+
| [1, 2, true, false]                   |
+---------------------------------------+
mysql> SELECT JSON_MERGE('{"name": "x"}', '{"id": 47}');
+-------------------------------------------+
| JSON_MERGE('{"name": "x"}', '{"id": 47}') |
+-------------------------------------------+
| {"id": 47, "name": "x"}                   |
+-------------------------------------------+
mysql> SELECT JSON_MERGE('1', 'true');
+------------------------+
| JSON_MERGE('1', 'true') |
+------------------------+
| [1, true]              |
+------------------------+
mysql> SELECT JSON_MERGE('[1, 2]', '{"id": 47}');
+-----------------------------------+
| JSON_MERGE('[1, 2]', '{"id": 47}') |
+-----------------------------------+
| [1, 2, {"id": 47}]                |
+-----------------------------------+
```

- `JSON_REMOVE(json_doc, path[, path] ...)`

Removes data from a JSON document and returns the result. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or is `$` or contains a `*` or `**` wildcard.

The *`path`* arguments are evaluated left to right. The document produced by evaluating one path becomes the new value against which the next path is evaluated.

It is not an error if the element to be removed does not exist in the document; in that case, the path does not affect the document.

```
mysql> SET @j = '["a", ["b", "c"], "d"]';
mysql> SELECT JSON_REMOVE(@j, '$[1]');
+-------------------------+
| JSON_REMOVE(@j, '$[1]') |
+-------------------------+
| ["a", "d"]              |
+-------------------------+
```

- `JSON_REPLACE(`*`json_doc`*`, `*`path`*`, `*`val`*`[, `*`path`*`, `*`val`*`] ...)`

  Replaces existing values in a JSON document and returns the result. Returns `NULL` if any argument is `NULL`. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

  A path-value pair for an existing path in the document overwrites the existing document value with the new value. A path-value pair for a nonexisting path in the document is ignored and has no effect.

  For a comparison of `JSON_INSERT()`, `JSON_REPLACE()`, and `JSON_SET()`, see the discussion of `JSON_SET()`.

```
mysql> SET @j = '{ "a": 1, "b": [2, 3]}';
mysql> SELECT JSON_REPLACE(@j, '$.a', 10, '$.c', '[true, false]');
+-----------------------------------------------------+
| JSON_REPLACE(@j, '$.a', 10, '$.c', '[true, false]') |
+-----------------------------------------------------+
| {"a": 10, "b": [2, 3]}                              |
+-----------------------------------------------------+
```

- `JSON_SET(`*`json_doc`*`, `*`path`*`, `*`val`*`[, `*`path`*`, `*`val`*`] ...)`

  Inserts or updates data in a JSON document and returns the result. Returns `NULL` if any argument is `NULL` or *`path`*, if given, does not locate an object. An error occurs if the *`json_doc`* argument is not a valid JSON document or any *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The path-value pairs are evaluated left to right. The document produced by evaluating one pair becomes the new value against which the next pair is evaluated.

  A path-value pair for an existing path in the document overwrites the existing document value with the new value. A path-value pair for a nonexisting path in the document adds the value to the document if the path identifies one of these types of values:

  - A member not present in an existing object. The member is added to the object and associated with the new value.
  - A position past the end of an existing array. The array is extended with the new value. If the existing value is not an array, it is autowrapped as an array, then extended with the new value.

Otherwise, a path-value pair for a nonexisting path in the document is ignored and has no effect.

The `JSON_SET()`, `JSON_INSERT()`, and `JSON_REPLACE()` functions are related:

- `JSON_SET()` replaces existing values and adds nonexisting values.
- `JSON_INSERT()` inserts values without replacing existing values.
- `JSON_REPLACE()` replaces *only* existing values.

The following examples illustrate these differences, using one path that does exist in the document (`$.a`) and another that does not exist (`$.c`):

```
mysql> SET @j = '{ "a": 1, "b": [2, 3]}';
mysql> SELECT JSON_SET(@j, '$.a', 10, '$.c', '[true, false]');
+-------------------------------------------------+
| JSON_SET(@j, '$.a', 10, '$.c', '[true, false]') |
+-------------------------------------------------+
| {"a": 10, "b": [2, 3], "c": "[true, false]"}    |
+-------------------------------------------------+
mysql> SELECT JSON_INSERT(@j, '$.a', 10, '$.c', '[true, false]');
+----------------------------------------------------+
| JSON_INSERT(@j, '$.a', 10, '$.c', '[true, false]') |
+----------------------------------------------------+
| {"a": 1, "b": [2, 3], "c": "[true, false]"}        |
+----------------------------------------------------+
mysql> SELECT JSON_REPLACE(@j, '$.a', 10, '$.c', '[true, false]');
+-----------------------------------------------------+
| JSON_REPLACE(@j, '$.a', 10, '$.c', '[true, false]') |
+-----------------------------------------------------+
| {"a": 10, "b": [2, 3]}                              |
+-----------------------------------------------------+
```

- `JSON_UNQUOTE(`*`json_val`*`)`

Unquotes JSON value and returns the result as a `utf8mb4` string. Returns `NULL` if the argument is `NULL`. An error occurs if the value starts and ends with double quotes but is not a valid JSON string literal.

Within a string, certain sequences have special meaning unless the `NO_BACKSLASH_ESCAPES` SQL mode is enabled. Each of these sequences begins with a backslash (`\`), known as the *escape character*. MySQL recognizes the escape sequences shown in Table 12.21, "JSON_UNQUOTE() Special Character Escape Sequences". For all other escape sequences, backslash is ignored. That is, the escaped character is interpreted as if it was not escaped. For example, `\x` is just `x`. These sequences are case sensitive. For example, `\b` is interpreted as a backspace, but `\B` is interpreted as `B`.

Table 12.21 JSON_UNQUOTE() Special Character Escape Sequences

Escape Sequence	Character Represented by Sequence
`\"`	A double quote (") character
`\b`	A backspace character
`\f`	A formfeed character
`\n`	A newline (linefeed) character
`\r`	A carriage return character
`\t`	A tab character
`\\`	A backslash (\) character
`\u`*`XXXX`*	UTF-8 bytes for Unicode value *`XXXX`*

Two simple examples of the use of this function are shown here:

```
mysql> SET @j = '"abc"';
mysql> SELECT @j, JSON_UNQUOTE(@j);
+-------+------------------+
| @j    | JSON_UNQUOTE(@j) |
+-------+------------------+
| "abc" | abc              |
+-------+------------------+
mysql> SET @j = '[1, 2, 3]';
mysql> SELECT @j, JSON_UNQUOTE(@j);
+-----------+------------------+
| @j        | JSON_UNQUOTE(@j) |
+-----------+------------------+
| [1, 2, 3] | [1, 2, 3]        |
+-----------+------------------+
```

The following set of examples shows how `JSON_UNQUOTE` handles escapes with `NO_BACKSLASH_ESCAPES` disabled and enabled:

```
mysql> SELECT @@sql_mode;
+------------+
| @@sql_mode |
+------------+
|            |
+------------+

mysql> SELECT JSON_UNQUOTE('"\\t\\u0032"');
+------------------------------+
| JSON_UNQUOTE('"\\t\\u0032"') |
+------------------------------+
|       2                          |
+------------------------------+

mysql> SET @@sql_mode = 'NO_BACKSLASH_ESCAPES';
mysql> SELECT JSON_UNQUOTE('"\\t\\u0032"');
+------------------------------+
| JSON_UNQUOTE('"\\t\\u0032"') |
+------------------------------+
| \t\u0032                     |
+------------------------------+

mysql> SELECT JSON_UNQUOTE('"\t\u0032"');
+---------------------------+
| JSON_UNQUOTE('"\t\u0032"') |
+---------------------------+
|       2                       |
+---------------------------+
```

12.16.5 Functions That Return JSON Value Attributes

The functions in this section return attributes of JSON values.

- `JSON_DEPTH(`*`json_doc`*`)`

  Returns the maximum depth of a JSON document. Returns `NULL` if the argument is `NULL`. An error occurs if the argument is not a valid JSON document.

  An empty array, empty object, or scalar value has depth 1. A nonempty array containing only elements of depth 1 or nonempty object containing only member values of depth 1 has depth 2. Otherwise, a JSON document has depth greater than 2.

```
mysql> SELECT JSON_DEPTH('{}'), JSON_DEPTH('[]'), JSON_DEPTH('true');
+-----------------+-----------------+-------------------+
| JSON_DEPTH('{}') | JSON_DEPTH('[]') | JSON_DEPTH('true') |
```

```
+-----------------+-----------------+-------------------+
|               1 |               1 |                 1 |
+-----------------+-----------------+-------------------+
mysql> SELECT JSON_DEPTH('[10, 20]'), JSON_DEPTH('[[], {}]');
+----------------------+----------------------+
| JSON_DEPTH('[10, 20]') | JSON_DEPTH('[[], {}]') |
+----------------------+----------------------+
|                    2 |                    2 |
+----------------------+----------------------+
mysql> SELECT JSON_DEPTH('[10, {"a": 20}]');
+------------------------------+
| JSON_DEPTH('[10, {"a": 20}]') |
+------------------------------+
|                            3 |
+------------------------------+
```

- `JSON_LENGTH(`*`json_doc`*`[, `*`path`*`])`

  Returns the length of a JSON document, or, if a *`path`* argument is given, the length of the value within the document identified by the path. Returns `NULL` if any argument is `NULL` or the *`path`* argument does not identify a value in the document. An error occurs if the *`json_doc`* argument is not a valid JSON document or the *`path`* argument is not a valid path expression or contains a `*` or `**` wildcard.

  The length of a document is determined as follows:

  - The length of a scalar is 1.
  - The length of an array is the number of array elements.
  - The length of an object is the number of object members.
  - The length does not count the length of nested arrays or objects.

```
mysql> SELECT JSON_LENGTH('[1, 2, {"a": 3}]');
+---------------------------------+
| JSON_LENGTH('[1, 2, {"a": 3}]') |
+---------------------------------+
|                               3 |
+---------------------------------+
mysql> SELECT JSON_LENGTH('{"a": 1, "b": {"c": 30}}');
+----------------------------------------+
| JSON_LENGTH('{"a": 1, "b": {"c": 30}}') |
+----------------------------------------+
|                                      2 |
+----------------------------------------+
mysql> SELECT JSON_LENGTH('{"a": 1, "b": {"c": 30}}', '$.b');
+-----------------------------------------------+
| JSON_LENGTH('{"a": 1, "b": {"c": 30}}', '$.b') |
+-----------------------------------------------+
|                                             1 |
+-----------------------------------------------+
```

- `JSON_TYPE(`*`json_val`*`)`

  Returns a `utf8mb4` string indicating the type of a JSON value. This can be an object, an array, or a scalar type, as shown here:

```
mysql> SET @j = '{"a": [10, true]}';
mysql> SELECT JSON_TYPE(@j);
+---------------+
| JSON_TYPE(@j) |
+---------------+
| OBJECT        |
+---------------+
mysql> SELECT JSON_TYPE(JSON_EXTRACT(@j, '$.a'));
+-----------------------------------+
```

```
| JSON_TYPE(JSON_EXTRACT(@j, '$.a')) |
+------------------------------------+
| ARRAY                              |
+------------------------------------+
mysql> SELECT JSON_TYPE(JSON_EXTRACT(@j, '$.a[0]'));
+---------------------------------------+
| JSON_TYPE(JSON_EXTRACT(@j, '$.a[0]')) |
+---------------------------------------+
| INTEGER                               |
+---------------------------------------+
mysql> SELECT JSON_TYPE(JSON_EXTRACT(@j, '$.a[1]'));
+---------------------------------------+
| JSON_TYPE(JSON_EXTRACT(@j, '$.a[1]')) |
+---------------------------------------+
| BOOLEAN                               |
+---------------------------------------+
```

`JSON_TYPE()` returns `NULL` if the argument is `NULL`:

```
mysql> SELECT JSON_TYPE(NULL);
+-----------------+
| JSON_TYPE(NULL) |
+-----------------+
| NULL            |
+-----------------+
```

An error occurs if the argument is not a valid JSON value:

```
mysql> SELECT JSON_TYPE(1);
ERROR 3146 (22032): Invalid data type for JSON data in argument 1
to function json_type; a JSON string or JSON type is required.
```

For a non-`NULL`, non-error result, the following list describes the possible `JSON_TYPE()` return values:

- Purely JSON types:
  - `OBJECT`: JSON objects
  - `ARRAY`: JSON arrays
  - `BOOLEAN`: The JSON true and false literals
  - `NULL`: The JSON null literal
- Numeric types:
  - `INTEGER`: MySQL `TINYINT`, `SMALLINT`, `MEDIUMINT` and `INT` and `BIGINT` scalars
  - `DOUBLE`: MySQL `DOUBLE` `FLOAT` scalars
  - `DECIMAL`: MySQL `DECIMAL` and `NUMERIC` scalars
- Temporal types:
  - `DATETIME`: MySQL `DATETIME` and `TIMESTAMP` scalars
  - `DATE`: MySQL `DATE` scalars
  - `TIME`: MySQL `TIME` scalars
- String types:
  - `STRING`: MySQL `utf8` character type scalars: `CHAR`, `VARCHAR`, `TEXT`, `ENUM`, and `SET`

  - Binary types:
    - `BLOB`: MySQL binary type scalars including `BINARY`, `VARBINARY`, `BLOB`, and `BIT`
  - All other types:
    - `OPAQUE` (raw bits)

- `JSON_VALID(val)`

  Returns 0 or 1 to indicate whether a value is valid JSON. Returns `NULL` if the argument is `NULL`.

```
mysql> SELECT JSON_VALID('{"a": 1}');
+------------------------+
| JSON_VALID('{"a": 1}') |
+------------------------+
|                      1 |
+------------------------+
mysql> SELECT JSON_VALID('hello'), JSON_VALID('"hello"');
+---------------------+-----------------------+
| JSON_VALID('hello') | JSON_VALID('"hello"') |
+---------------------+-----------------------+
|                   0 |                     1 |
+---------------------+-----------------------+
```

12.16.6 JSON Utility Functions

This section documents utility functions that act on JSON values, or strings that can be parsed as JSON values. `JSON_PRETTY()` prints out a JSON value in a format that is easy to read. `JSON_STORAGE_SIZE()` and `JSON_STORAGE_FREE()` show, respectively, the amount of storage space used by a given JSON value and the amount of space remaining in a `JSON` column following a partial update.

- `JSON_PRETTY(json_val)`

  Provides pretty-printing of JSON values similar to that implemented in PHP and by other languages and database systems. The value supplied must be a JSON value or a valid string representation of a JSON value. Extraneous whitespaces and newlines present in this value have no effect on the output. For a `NULL` value, the function returns `NULL`. If the value is not a JSON document, or if it cannot cannot be parsed as one, the function fails with an error.

  Formatting of the output from this function adheres to the following rules:

  - Each array element or object member appears on a separate line, indented by one additional level as compared to its parent.
  - Each level of indentation adds two leading spaces.
  - A comma separating individual array elements or object members is printed before the newline that separates the two elements or members.
  - The key and the value of an object member are separated by a colon followed by a space (`': '`).
  - An empty object or array is printed on a single line. No space is printed between the opening and closing brace.
  - Special characters in string scalars and key names are escaped employing the same rules used by the `JSON_QUOTE()` function.

```
mysql> SELECT JSON_PRETTY('123'); # scalar
+--------------------+
| JSON_PRETTY('123') |
```

```
+--------------------+
| 123                |
+--------------------+

mysql> SELECT JSON_PRETTY("[1,3,5]"); # array
+------------------------+
| JSON_PRETTY("[1,3,5]") |
+------------------------+
| [
  1,
  3,
  5
]      |
+------------------------+

mysql> SELECT JSON_PRETTY('{"a":"10","b":"15","x":"25"}'); # object
+-------------------------------------------+
| JSON_PRETTY('{"a":"10","b":"15","x":"25"}') |
+-------------------------------------------+
| {
  "a": "10",
  "b": "15",
  "x": "25"
}   |
+-------------------------------------------+

mysql> SELECT JSON_PRETTY('["a",1,{"key1":
     >    "value1"},"5",     "77" ,
     >       {"key2":["value3","valueX",
     > "valueY"]},"j", "2"   ]')\G  # nested arrays and objects
*************************** 1. row ***************************
JSON_PRETTY('["a",1,{"key1":
             "value1"},"5",     "77" ,
                {"key2":["value3","valuex",
          "valuey"]},"j", "2"   ]'): [
  "a",
  1,
  {
    "key1": "value1"
  },
  "5",
  "77",
  {
    "key2": [
      "value3",
      "valuex",
      "valuey"
    ]
  },
  "j",
  "2"
]
```

- `JSON_STORAGE_FREE(json_val)`

  For a `JSON` column value, this function shows how much storage space was freed in its binary representation after it was updated in place using `JSON_SET()`, `JSON_REPLACE()`, or `JSON_REMOVE()`. The argument can also be a valid JSON document or a string which can be parsed as one—either as a literal value or as the value of a user variable—in which case the function returns 0. It returns a positive, nonzero value if the argument is a `JSON` column value which has been updated as described previously, such that its binary representation takes up less space than it did prior to the update. For a `JSON` column which has been updated such that its binary representation is the same as or larger than before, or if the update was not able to take advantage of a partial update, it returns 0; it returns `NULL` if the argument is `NULL`.

  If *`json_val`* is not `NULL`, and neither is a valid JSON document nor can be successfully parsed as one, an error results.

In this example, we create a table containing a `JSON` column, then insert a row containing a JSON object:

```
mysql> CREATE TABLE jtable (jcol JSON);
Query OK, 0 rows affected (0.38 sec)

mysql> INSERT INTO jtable VALUES
    ->     ('{"a": 10, "b": "wxyz", "c": "[true, false]"}');
Query OK, 1 row affected (0.04 sec)

mysql> SELECT * FROM jtable;
+-------------------------------------------+
| jcol                                      |
+-------------------------------------------+
| {"a": 10, "b": "wxyz", "c": "[true, false]"} |
+-------------------------------------------+
1 row in set (0.00 sec)
```

Now we update the column value using `JSON_SET()` such that a partial update can be performed; in this case, we replace the value pointed to by the `c` key (the array `[true, false]`) with one that takes up less space (the integer `1`):

```
mysql> UPDATE jtable
    ->     SET jcol = JSON_SET(jcol, "$.a", 10, "$.b", "wxyz", "$.c", 1);
Query OK, 1 row affected (0.03 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT * FROM jtable;
+-------------------------------+
| jcol                          |
+-------------------------------+
| {"a": 10, "b": "wxyz", "c": 1} |
+-------------------------------+
1 row in set (0.00 sec)

mysql> SELECT JSON_STORAGE_FREE(jcol) FROM jtable;
+-------------------------+
| JSON_STORAGE_FREE(jcol) |
+-------------------------+
|                      14 |
+-------------------------+
1 row in set (0.00 sec)
```

Updating the column without using `JSON_SET()`, `JSON_REPLACE()`, or `JSON_REMOVE()` means that the optimizer cannot perform the update in place; in this case, `JSON_STORAGE_FREE()` returns 0, as shown here:

```
mysql> UPDATE jtable SET jcol = '{"a": 10, "b": 1}';
Query OK, 1 row affected (0.05 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT JSON_STORAGE_FREE(jcol) FROM jtable;
+-------------------------+
| JSON_STORAGE_FREE(jcol) |
+-------------------------+
|                       0 |
+-------------------------+
1 row in set (0.00 sec)
```

Partial updates of JSON documents can be performed only on column values. For a user variable that stores a JSON value, the value is always completely replaced, even when the update is performed using `JSON_SET()`:

```
mysql> SET @j = '{"a": 10, "b": "wxyz", "c": "[true, false]"}';
```

```
Query OK, 0 rows affected (0.00 sec)

mysql> SET @j = JSON_SET(@j, '$.a', 10, '$.b', 'wxyz', '$.c', '1');
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @j, JSON_STORAGE_FREE(@j) AS Free;
+---------------------------------+------+
| @j                              | Free |
+---------------------------------+------+
| {"a": 10, "b": "wxyz", "c": "1"} |    0 |
+---------------------------------+------+
1 row in set (0.00 sec)
```

For a JSON literal, this function always returns 0:

```
mysql> SELECT JSON_STORAGE_FREE('{"a": 10, "b": "wxyz", "c": "1"}') AS Free;
+------+
| Free |
+------+
|    0 |
+------+
1 row in set (0.00 sec)
```

- `JSON_STORAGE_SIZE(`*`json_val`*`)`

  This function returns the number of bytes used to store the binary representation of a JSON document. When the argument is a `JSON` column, this is the space used to store the JSON document as it was inserted into the column, prior to any partial updates that may have been performed on it afterwards. *`json_val`* must be a valid JSON document or a string which can be parsed as one. In the case where it is string, the function returns the amount of storage space in the JSON binary representation that is created by parsing the string as JSON and converting it to binary. It returns `NULL` if the argument is `NULL`.

  An error results when *`json_val`* is not `NULL`, and is not—or cannot be successfully parsed as—a JSON document.

  To illustrate this function's behavior when used with a `JSON` column as its argument, we create a table named `jtable` containing a `JSON` column `jcol`, insert a JSON value into the table, then obtain the storage space used by this column with `JSON_STORAGE_SIZE()`, as shown here:

```
mysql> CREATE TABLE jtable (jcol JSON);
Query OK, 0 rows affected (0.42 sec)

mysql> INSERT INTO jtable VALUES
    ->     ('{"a": 1000, "b": "wxyz", "c": "[1, 3, 5, 7]"}');
Query OK, 1 row affected (0.04 sec)

mysql> SELECT
    ->     jcol,
    ->     JSON_STORAGE_SIZE(jcol) AS Size,
    ->     JSON_STORAGE_FREE(jcol) AS Free
    -> FROM jtable;
+-----------------------------------------------+------+------+
| jcol                                          | Size | Free |
+-----------------------------------------------+------+------+
| {"a": 1000, "b": "wxyz", "c": "[1, 3, 5, 7]"} |   47 |    0 |
+-----------------------------------------------+------+------+
1 row in set (0.00 sec)
```

  According to the output of `JSON_STORAGE_SIZE()`, the JSON document inserted into the column takes up 47 bytes. We also checked the amount of space freed by any previous partial updates of the column using `JSON_STORAGE_FREE()`; since no updates have yet been performed, this is 0, as expected.

Next we perform an `UPDATE` on the table that should result in a partial update of the document stored in `jcol`, and then test the result as shown here:

```
mysql> UPDATE jtable SET jcol =
    ->     JSON_SET(jcol, "$.b", "a");
Query OK, 1 row affected (0.04 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT
    ->     jcol,
    ->     JSON_STORAGE_SIZE(jcol) AS Size,
    ->     JSON_STORAGE_FREE(jcol) AS Free
    -> FROM jtable;
+-----------------------------------------+------+------+
| jcol                                    | Size | Free |
+-----------------------------------------+------+------+
| {"a": 1000, "b": "a", "c": "[1, 3, 5, 7]"} |   47 |    3 |
+-----------------------------------------+------+------+
1 row in set (0.00 sec)
```

The value returned by `JSON_STORAGE_FREE()` in the previous query indicates that a partial update of the JSON document was performed, and that this freed 3 bytes of space used to store it. The result returned by `JSON_STORAGE_SIZE()` is unchanged by the partial update.

Partial updates are supported for updates using `JSON_SET()`, `JSON_REPLACE()`, or `JSON_REMOVE()`. The direct assignment of a value to a `JSON` column cannot be partially updated; following such an update, `JSON_STORAGE_SIZE()` always shows the storage used for the newly-set value:

```
mysql> UPDATE jtable
mysql>     SET jcol = '{"a": 4.55, "b": "wxyz", "c": "[true, false]"}';
Query OK, 1 row affected (0.04 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT
    ->     jcol,
    ->     JSON_STORAGE_SIZE(jcol) AS Size,
    ->     JSON_STORAGE_FREE(jcol) AS Free
    -> FROM jtable;
+-------------------------------------------------+------+------+
| jcol                                            | Size | Free |
+-------------------------------------------------+------+------+
| {"a": 4.55, "b": "wxyz", "c": "[true, false]"} |   56 |    0 |
+-------------------------------------------------+------+------+
1 row in set (0.00 sec)
```

A JSON user variable cannot be partially updated. This means that this function always shows the space currently used to store a JSON document in a user variable:

```
mysql> SET @j = '[100, "sakila", [1, 3, 5], 425.05]';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @j, JSON_STORAGE_SIZE(@j) AS Size;
+-----------------------------------+------+
| @j                                | Size |
+-----------------------------------+------+
| [100, "sakila", [1, 3, 5], 425.05] |   45 |
+-----------------------------------+------+
1 row in set (0.00 sec)

mysql> SET @j = JSON_SET(@j, '$[1]', "json");
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @j, JSON_STORAGE_SIZE(@j) AS Size;
+-----------------------------------+------+
```

```
| @j                                | Size |
+-----------------------------------+------+
| [100, "json", [1, 3, 5], 425.05] |   43 |
+-----------------------------------+------+
1 row in set (0.00 sec)

mysql> SET @j = JSON_SET(@j, '$[2][0]', JSON_ARRAY(10, 20, 30));
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @j, JSON_STORAGE_SIZE(@j) AS Size;
+-------------------------------------------+------+
| @j                                        | Size |
+-------------------------------------------+------+
| [100, "json", [[10, 20, 30], 3, 5], 425.05] |   56 |
+-------------------------------------------+------+
1 row in set (0.00 sec)
```

For a JSON literal, this function always returns the current storage space used:

```
mysql> SELECT
    ->     JSON_STORAGE_SIZE('[100, "sakila", [1, 3, 5], 425.05]') AS A,
    ->     JSON_STORAGE_SIZE('{"a": 1000, "b": "a", "c": "[1, 3, 5, 7]"}') AS B,
    ->     JSON_STORAGE_SIZE('{"a": 1000, "b": "wxyz", "c": "[1, 3, 5, 7]"}') AS C,
    ->     JSON_STORAGE_SIZE('[100, "json", [[10, 20, 30], 3, 5], 425.05]') AS D;
+----+----+----+----+
| A  | B  | C  | D  |
+----+----+----+----+
| 45 | 44 | 47 | 56 |
+----+----+----+----+
1 row in set (0.00 sec)
```

12.16.7 JSON Path Syntax

Many of the functions described in previous sections require a path expression in order to identify a specific element in a JSON document. A path consists of the path's scope followed by one or more path legs. For paths used in MySQL JSON functions, the scope is always the document being searched or otherwise operated on, represented by a leading `$` character. Path legs are separated by period characters (`.`). Cells in arrays are represented by `[N]`, where *`N`* is a non-negative integer. Names of keys must be double-quoted strings or valid ECMAScript identifiers (see `http://www.ecma-international.org/ecma-262/5.1/#sec-7.6`). Path expressions, like JSON text, should be encoded using the `ascii`, `utf8`, or `utf8mb4` character sets. Other character encodings are implicitly coerced to `utf8mb4`. The complete syntax is shown here:

```
pathExpression:
    scope[(pathLeg)*]

pathLeg:
    member | arrayLocation | doubleAsterisk

member:
    period ( keyName | asterisk )

arrayLocation:
    leftBracket ( nonNegativeInteger | asterisk ) rightBracket

keyName:
    ESIdentifier | doubleQuotedString

doubleAsterisk:
    '**'

period:
    '.'

asterisk:
    '*'
```

```
leftBracket:
    '['

rightBracket:
    ']'
```

As noted previously, in MySQL, the scope of the path is always the document being operated on, represented as `$`. You can use `'$'` as a synonym for the document in JSON path expressions.

Note

Some implementations support column references for scopes of JSON paths; currently, MySQL does not support these.

The wildcard `*` and `**` tokens are used as follows:

- `.*` represents the values of all members in the object.
- `[*]` represents the values of all cells in the array.
- `[`*`prefix`*`]**`*`suffix`* represents all paths beginning with *`prefix`* and ending with *`suffix`*. *`prefix`* is optional, while *`suffix`* is required; in other words, a path may not end in `**`.

  In addition, a path may not contain the sequence `***`.

For path syntax examples, see the descriptions of the various JSON fuinctions that take paths as arguments, such as `JSON_CONTAINS_PATH()`, `JSON_SET()`, and `JSON_REPLACE()`. For examples which include the use of the `*` and `**` wildcards, see the description of the `JSON_SEARCH()` function.

12.17 Functions Used with Global Transaction IDs

The functions described in this section are used with GTID-based replication. It is important to keep in mind that all of these functions take string representations of GTID sets as arguments—as such, the GTID sets must always be quoted when used with them. See GTID Sets for more information.

The union of two GTID sets is simply their representations as strings, joined together with an interposed comma. In other words, you can define a very simple function for obtaining the union of two GTID sets, similar to that created here:

```
CREATE FUNCTION GTID_UNION(g1 TEXT, g2 TEXT)
    RETURNS TEXT DETERMINISTIC
    RETURN CONCAT(g1,',',g2);
```

For more information about GTIDs and how these GTID functions are used in practice, see Section 18.1.3, “Replication with Global Transaction Identifiers”.

Table 12.22 GTID Functions

Name	Description
`GTID_SUBSET()`	Return true if all GTIDs in subset are also in set; otherwise false.
`GTID_SUBTRACT()`	Return all GTIDs in set that are not in subset.
`WAIT_FOR_EXECUTED_GTID_SET()`	Wait until the given GTIDs have executed on slave.
`WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()`	Wait until the given GTIDs have executed on slave.

- `GTID_SUBSET(`*`subset`*`,`*`set`*`)`

  Given two sets of global transaction IDs *`subset`* and *`set`*, returns true if all GTIDs in *`subset`* are also in *`set`*. Returns false otherwise.

  The GTID sets used with this function are represented as strings, as shown in the following examples:

```
mysql> SELECT GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:23',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57')\G
*************************** 1. row ***************************
GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:23',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57'): 1
1 row in set (0.00 sec)

mysql> SELECT GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:23-25',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57')\G
*************************** 1. row ***************************
GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:23-25',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57'): 1
1 row in set (0.00 sec)

mysql> SELECT GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:20-25',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57')\G
*************************** 1. row ***************************
GTID_SUBSET('3E11FA47-71CA-11E1-9E33-C80AA9429562:20-25',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57'): 0
1 row in set (0.00 sec)
```

- `GTID_SUBTRACT(set,subset)`

  Given two sets of global transaction IDs *`subset`* and *`set`*, returns only those GTIDs from *`set`* that are not in *`subset`*.

  All GTID sets used with this function are represented as strings and must be quoted, as shown in these examples:

```
mysql> SELECT GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:21')\G
*************************** 1. row ***************************
GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:21'): 3e11fa47-71ca-11e1-9e33-c80aa9429562:22-57
1 row in set (0.00 sec)

mysql> SELECT GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:20-25')\G
*************************** 1. row ***************************
GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:20-25'): 3e11fa47-71ca-11e1-9e33-c80aa9429562:26-57
1 row in set (0.00 sec)

mysql> SELECT GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    ->     '3E11FA47-71CA-11E1-9E33-C80AA9429562:23-24')\G
*************************** 1. row ***************************
GTID_SUBTRACT('3E11FA47-71CA-11E1-9E33-C80AA9429562:21-57',
    '3E11FA47-71CA-11E1-9E33-C80AA9429562:23-24'): 3e11fa47-71ca-11e1-9e33-c80aa9429562:21-22:25-57
1 row in set (0.01 sec)
```

- `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS(gtid_set[, timeout][,channel])`

  Wait until the slave SQL thread has executed all of the transactions whose global transaction identifiers are contained in *`gtid_set`* (see Section 18.1.3.1, “GTID Concepts”, for a definition of “GTID sets”), or until *`timeout`* seconds have elapsed, whichever occurs first. *`timeout`* is optional; the default timeout is 0 seconds, in which case the function waits until all of the transactions in the GTID set have been executed.

  For more information, see Section 18.1.3, “Replication with Global Transaction Identifiers”.

  GTID sets used with this function are represented as strings and so must be quoted as shown in the following example:

```
mysql> SELECT WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS('3E11FA47-71CA-11E1-9E33-C80AA9429562:1-5');
        -> 5
```

The return value is the number of transactional events that were executed. If GTID-based replication is not active (that is, if the value of the `gtid_mode` variable is `OFF`), then this value is undefined and thus `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()` returns `NULL`. If the slave is not running then this function also returns `NULL`.

The *`channel`* option enables you to choose which replication channel the function applies to. If no *`channel`* is set and no channels other than the default replication channel exist, the function applies to the default replication channel. If you are using multiple replication channels you must specify a *`channel`* as otherwise it is not known which replication channel the function should act on. See Section 18.2.3, "Replication Channels" for more information on replication channels.

- `WAIT_FOR_EXECUTED_GTID_SET(`*`gtid_set`*`[, `*`timeout`*`])`

  `WAIT_FOR_EXECUTED_GTID_SET()` is similar to `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()` in that it waits until a server has executed all of the transactions whose global transaction identifiers are contained in *`gtid_set`*, or until *`timeout`* seconds have elapsed, whichever occurs first. Unlike `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()`, `WAIT_FOR_EXECUTED_GTID_SET()` does not take into account whether the slave is running or not, and an error is returned if GTID-based replication is not enabled.

  In addition, `WAIT_FOR_EXECUTED_GTID_SET()` returns only the state of the query, where 0 represents success, 1 represents timeout, and any other failures return the error message.

12.18 Aggregate (GROUP BY) Functions

12.18.1 Aggregate (GROUP BY) Function Descriptions

Table 12.23 Aggregate (`GROUP BY`) Functions

Name	Description
`AVG()`	Return the average value of the argument
`BIT_AND()`	Return bitwise AND
`BIT_OR()`	Return bitwise OR
`BIT_XOR()`	Return bitwise XOR
`COUNT()`	Return a count of the number of rows returned
`COUNT(DISTINCT)`	Return the count of a number of different values
`GROUP_CONCAT()`	Return a concatenated string
`JSON_ARRAYAGG()`	Return result set as a single JSON array
`JSON_OBJECTAGG()`	Return result set as a single JSON object
`MAX()`	Return the maximum value
`MIN()`	Return the minimum value
`STD()`	Return the population standard deviation
`STDDEV()`	Return the population standard deviation
`STDDEV_POP()`	Return the population standard deviation
`STDDEV_SAMP()`	Return the sample standard deviation
`SUM()`	Return the sum
`VAR_POP()`	Return the population standard variance
`VAR_SAMP()`	Return the sample variance
`VARIANCE()`	Return the population standard variance

This section describes group (aggregate) functions that operate on sets of values. Unless otherwise stated, group functions ignore `NULL` values.

If you use a group function in a statement containing no `GROUP BY` clause, it is equivalent to grouping on all rows. For more information, see Section 12.18.3, “MySQL Handling of GROUP BY”.

For numeric arguments, the variance and standard deviation functions return a `DOUBLE` value. The `SUM()` and `AVG()` functions return a `DECIMAL` value for exact-value arguments (integer or `DECIMAL`), and a `DOUBLE` value for approximate-value arguments (`FLOAT` or `DOUBLE`).

The `SUM()` and `AVG()` aggregate functions do not work with temporal values. (They convert the values to numbers, losing everything after the first nonnumeric character.) To work around this problem, convert to numeric units, perform the aggregate operation, and convert back to a temporal value. Examples:

```
SELECT SEC_TO_TIME(SUM(TIME_TO_SEC(time_col))) FROM tbl_name;
SELECT FROM_DAYS(SUM(TO_DAYS(date_col))) FROM tbl_name;
```

Functions such as `SUM()` or `AVG()` that expect a numeric argument cast the argument to a number if necessary. For `SET` or `ENUM` values, the cast operation causes the underlying numeric value to be used.

The `BIT_AND()`, `BIT_OR()`, and `BIT_XOR()` aggregate functions perform bit operations. Prior to MySQL 8.0, bit functions and operators required `BIGINT` (64-bit integer) arguments and returned `BIGINT` values, so they had a maximum range of 64 bits. Non-`BIGINT` arguments were converted to `BIGINT` prior to performing the operation and truncation could occur.

In MySQL 8.0, bit functions and operators permit binary string type arguments (`BINARY`, `VARBINARY`, and the `BLOB` types) and return a value of like type, which enables them to take arguments and produce return values larger than 64 bits. For discussion about argument evaluation and result types for bit operations, see the introductory discussion in Section 12.12, “Bit Functions and Operators”.

- `AVG([DISTINCT] expr)`

  Returns the average value of *`expr`*. The `DISTINCT` option can be used to return the average of the distinct values of *`expr`*.

  If there are no matching rows, `AVG()` returns `NULL`.

  ```
  mysql> SELECT student_name, AVG(test_score)
      ->        FROM student
      ->        GROUP BY student_name;
  ```

- `BIT_AND(expr)`

  Returns the bitwise `AND` of all bits in *`expr`*.

  The result type depends on whether the function argument values are evaluated as binary strings or numbers:

  - Binary-string evaluation occurs when the argument values have a binary string type, and the argument is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument value conversion to unsigned 64-bit integers as necessary.

  - Binary-string evaluation produces a binary string of the same length as the argument values. If argument values have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. If the argument size exceeds 511 bytes, an `ER_INVALID_BITWISE_AGGREGATE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

If there are no matching rows, `BIT_AND()` returns a neutral value (all bits set to 1) having the same length as the argument values.

`NULL` values do not affect the result unless all values are `NULL`. In that case, the result is a neutral value having the same length as the argument values.

For more information discussion about argument evaluation and result types, see the introductory discussion in Section 12.12, “Bit Functions and Operators”.

- `BIT_OR(`*`expr`*`)`

  Returns the bitwise `OR` of all bits in *`expr`*.

  The result type depends on whether the function argument values are evaluated as binary strings or numbers:

  - Binary-string evaluation occurs when the argument values have a binary string type, and the argument is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument value conversion to unsigned 64-bit integers as necessary.

  - Binary-string evaluation produces a binary string of the same length as the argument values. If argument values have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. If the argument size exceeds 511 bytes, an `ER_INVALID_BITWISE_AGGREGATE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

  If there are no matching rows, `BIT_OR()` returns a neutral value (all bits set to 0) having the same length as the argument values.

  `NULL` values do not affect the result unless all values are `NULL`. In that case, the result is a neutral value having the same length as the argument values.

  For more information discussion about argument evaluation and result types, see the introductory discussion in Section 12.12, “Bit Functions and Operators”.

- `BIT_XOR(`*`expr`*`)`

  Returns the bitwise `XOR` of all bits in *`expr`*.

  The result type depends on whether the function argument values are evaluated as binary strings or numbers:

  - Binary-string evaluation occurs when the argument values have a binary string type, and the argument is not a hexadecimal literal, bit literal, or `NULL` literal. Numeric evaluation occurs otherwise, with argument value conversion to unsigned 64-bit integers as necessary.

  - Binary-string evaluation produces a binary string of the same length as the argument values. If argument values have unequal lengths, an `ER_INVALID_BITWISE_OPERANDS_SIZE` error occurs. If the argument size exceeds 511 bytes, an `ER_INVALID_BITWISE_AGGREGATE_OPERANDS_SIZE` error occurs. Numeric evaluation produces an unsigned 64-bit integer.

  If there are no matching rows, `BIT_XOR()` returns a neutral value (all bits set to 0) having the same length as the argument values.

  `NULL` values do not affect the result unless all values are `NULL`. In that case, the result is a neutral value having the same length as the argument values.

  For more information discussion about argument evaluation and result types, see the introductory discussion in Section 12.12, “Bit Functions and Operators”.

- `COUNT(expr)`

  Returns a count of the number of non-`NULL` values of *`expr`* in the rows retrieved by a `SELECT` statement. The result is a `BIGINT` value.

  If there are no matching rows, `COUNT()` returns `0`.

```
mysql> SELECT student.student_name,COUNT(*)
    ->        FROM student,course
    ->        WHERE student.student_id=course.student_id
    ->        GROUP BY student_name;
```

  `COUNT(*)` is somewhat different in that it returns a count of the number of rows retrieved, whether or not they contain `NULL` values.

  For `MyISAM` tables, `COUNT(*)` is optimized to return very quickly if the `SELECT` retrieves from one table, no other columns are retrieved, and there is no `WHERE` clause. For example:

```
mysql> SELECT COUNT(*) FROM student;
```

  This optimization only applies to `MyISAM` tables, because an exact row count is stored for this storage engine and can be accessed very quickly. `COUNT(1)` is only subject to the same optimization if the first column is defined as `NOT NULL`.

  For transactional storage engines such as `InnoDB`, storing an exact row count is problematic because multiple transactions may be occurring, each of which may affect the count. For more information about how `InnoDB` handles `COUNT(*)` operations, see Section 15.8.1.7, “Limits on InnoDB Tables”.

- `COUNT(DISTINCT expr,[expr...])`

  Returns a count of the number of rows with different non-`NULL` *`expr`* values.

  If there are no matching rows, `COUNT(DISTINCT)` returns `0`.

```
mysql> SELECT COUNT(DISTINCT results) FROM student;
```

  In MySQL, you can obtain the number of distinct expression combinations that do not contain `NULL` by giving a list of expressions. In standard SQL, you would have to do a concatenation of all expressions inside `COUNT(DISTINCT ...)`.

- `GROUP_CONCAT(expr)`

  This function returns a string result with the concatenated non-`NULL` values from a group. It returns `NULL` if there are no non-`NULL` values. The full syntax is as follows:

```
GROUP_CONCAT([DISTINCT] expr [,expr ...]
             [ORDER BY {unsigned_integer | col_name | expr}
                 [ASC | DESC] [,col_name ...]]
             [SEPARATOR str_val])
```

```
mysql> SELECT student_name,
    ->     GROUP_CONCAT(test_score)
    ->     FROM student
    ->     GROUP BY student_name;
```

  Or:

```
mysql> SELECT student_name,
    ->     GROUP_CONCAT(DISTINCT test_score
```

```
    ->               ORDER BY test_score DESC SEPARATOR ' ')
    ->     FROM student
    ->     GROUP BY student_name;
```

In MySQL, you can get the concatenated values of expression combinations. To eliminate duplicate values, use the `DISTINCT` clause. To sort values in the result, use the `ORDER BY` clause. To sort in reverse order, add the `DESC` (descending) keyword to the name of the column you are sorting by in the `ORDER BY` clause. The default is ascending order; this may be specified explicitly using the `ASC` keyword. The default separator between values in a group is comma (`,`). To specify a separator explicitly, use `SEPARATOR` followed by the string literal value that should be inserted between group values. To eliminate the separator altogether, specify `SEPARATOR ''`.

The result is truncated to the maximum length that is given by the `group_concat_max_len` system variable, which has a default value of 1024. The value can be set higher, although the effective maximum length of the return value is constrained by the value of `max_allowed_packet`. The syntax to change the value of `group_concat_max_len` at runtime is as follows, where *`val`* is an unsigned integer:

```
SET [GLOBAL | SESSION] group_concat_max_len = val;
```

The return value is a nonbinary or binary string, depending on whether the arguments are nonbinary or binary strings. The result type is `TEXT` or `BLOB` unless `group_concat_max_len` is less than or equal to 512, in which case the result type is `VARCHAR` or `VARBINARY`.

See also `CONCAT()` and `CONCAT_WS()`: Section 12.5, “String Functions”.

- `JSON_ARRAYAGG(col_or_expr)`

  Aggregates a result set as a single `JSON` array whose elements consist of the rows. The order of elements in this array is undefined. The function acts on a column or an expression that evaluates to a single value. Returns `NULL` if the result contains no rows, or in the event of an error.

```
mysql> SELECT col FROM t1;
+--------------------------------------+
| col                                  |
+--------------------------------------+
| {"key1": "value1", "key2": "value2"} |
| {"keyA": "valueA", "keyB": "valueB"} |
+--------------------------------------+
2 rows in set (0.00 sec)

mysql> SELECT JSON_ARRAYAGG(col) FROM t1;
+--------------------------------------------------------------------------------+
| JSON_ARRAYAGG(col)                                                             |
+--------------------------------------------------------------------------------+
| [{"key1": "value1", "key2": "value2"}, {"keyA": "valueA", "keyB": "valueB"}] |
+--------------------------------------------------------------------------------+
1 row in set (0.00 sec)
```

- `JSON_OBJECTAGG(key, value)`

  Takes two column names or expressions as arguments, the first of these being used as a key and the second as a value, and returns a JSON object containing key/value pairs. Elements with duplicate keys are skipped. (You can specify a constant as the key, but the result always contains at most one row element.) Returns `NULL` if the result contains no rows, or in the event of an error. An error occurs if any key name is `NULL` or the number of arguments is not equal to 2.

```
mysql> SELECT id, name FROM t3 WHERE id < 10;
+------+------+
| id   | name |
+------+------+
|    2 | joe  |
|    5 | fred |
```

```
+------+------+
2 rows in set (0.00 sec)

mysql> SELECT JSON_OBJECTAGG(id, name) FROM t3 WHERE id < 10;
+---------------------------+
| JSON_OBJECTAGG(id, name)  |
+---------------------------+
| {"2": "joe", "5": "fred"} |
+---------------------------+
1 row in set (0.00 sec)
```

- `MAX([DISTINCT] expr)`

  Returns the maximum value of *`expr`*. `MAX()` may take a string argument; in such cases, it returns the maximum string value. See Section 8.3.1, “How MySQL Uses Indexes”. The `DISTINCT` keyword can be used to find the maximum of the distinct values of *`expr`*, however, this produces the same result as omitting `DISTINCT`.

  If there are no matching rows, `MAX()` returns `NULL`.

```
mysql> SELECT student_name, MIN(test_score), MAX(test_score)
    ->        FROM student
    ->        GROUP BY student_name;
```

  For `MAX()`, MySQL currently compares `ENUM` and `SET` columns by their string value rather than by the string's relative position in the set. This differs from how `ORDER BY` compares them. This is expected to be rectified in a future MySQL release.

- `MIN([DISTINCT] expr)`

  Returns the minimum value of *`expr`*. `MIN()` may take a string argument; in such cases, it returns the minimum string value. See Section 8.3.1, “How MySQL Uses Indexes”. The `DISTINCT` keyword can be used to find the minimum of the distinct values of *`expr`*, however, this produces the same result as omitting `DISTINCT`.

  If there are no matching rows, `MIN()` returns `NULL`.

```
mysql> SELECT student_name, MIN(test_score), MAX(test_score)
    ->        FROM student
    ->        GROUP BY student_name;
```

  For `MIN()`, MySQL currently compares `ENUM` and `SET` columns by their string value rather than by the string's relative position in the set. This differs from how `ORDER BY` compares them. This is expected to be rectified in a future MySQL release.

- `STD(expr)`

  Returns the population standard deviation of *`expr`*. This is an extension to standard SQL. The standard SQL function `STDDEV_POP()` can be used instead.

  If there are no matching rows, `STD()` returns `NULL`.

- `STDDEV(expr)`

  Returns the population standard deviation of *`expr`*. This function is provided for compatibility with Oracle. The standard SQL function `STDDEV_POP()` can be used instead.

  If there are no matching rows, `STDDEV()` returns `NULL`.

- `STDDEV_POP(expr)`

  Returns the population standard deviation of *`expr`* (the square root of `VAR_POP()`). You can also use `STD()` or `STDDEV()`, which are equivalent but not standard SQL.

If there are no matching rows, `STDDEV_POP()` returns `NULL`.

- `STDDEV_SAMP(expr)`

  Returns the sample standard deviation of *`expr`* (the square root of `VAR_SAMP()`.

  If there are no matching rows, `STDDEV_SAMP()` returns `NULL`.

- `SUM([DISTINCT] expr)`

  Returns the sum of *`expr`*. If the return set has no rows, `SUM()` returns `NULL`. The `DISTINCT` keyword can be used to sum only the distinct values of *`expr`*.

  If there are no matching rows, `SUM()` returns `NULL`.

- `VAR_POP(expr)`

  Returns the population standard variance of *`expr`*. It considers rows as the whole population, not as a sample, so it has the number of rows as the denominator. You can also use `VARIANCE()`, which is equivalent but is not standard SQL.

  If there are no matching rows, `VAR_POP()` returns `NULL`.

- `VAR_SAMP(expr)`

  Returns the sample variance of *`expr`*. That is, the denominator is the number of rows minus one.

  If there are no matching rows, `VAR_SAMP()` returns `NULL`.

- `VARIANCE(expr)`

  Returns the population standard variance of *`expr`*. This is an extension to standard SQL. The standard SQL function `VAR_POP()` can be used instead.

  If there are no matching rows, `VARIANCE()` returns `NULL`.

12.18.2 GROUP BY Modifiers

The `GROUP BY` clause permits a `WITH ROLLUP` modifier that causes summary output to include extra rows that represent higher-level (that is, super-aggregate) summary operations. `ROLLUP` thus enables you to answer questions at multiple levels of analysis with a single query. For example, `ROLLUP` can be used to provide support for OLAP (Online Analytical Processing) operations.

Suppose that a `sales` table has `year`, `country`, `product`, and `profit` columns for recording sales profitability:

```
CREATE TABLE sales
(
    year    INT,
    country VARCHAR(20),
    product VARCHAR(32),
    profit  INT
);
```

To summarize table contents per year, use a simple `GROUP BY` like this:

```
mysql> SELECT year, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC;
+------+--------+
| year | profit |
```

```
+------+--------+
| 2000 |   4525 |
| 2001 |   3010 |
+------+--------+
```

The output shows the total profit for each year. To also determine the total profit summed over all years, you must add up the individual values yourself or run an additional query. Or you can use `ROLLUP`, which provides both levels of analysis with a single query. Adding a `WITH ROLLUP` modifier to the `GROUP BY` clause causes the query to produce another row that shows the grand total over all year values:

```
mysql> SELECT year, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC WITH ROLLUP;
+------+--------+
| year | profit |
+------+--------+
| 2000 |   4525 |
| 2001 |   3010 |
| NULL |   7535 |
+------+--------+
```

The `NULL` value in the `year` column identifies the grand total super-aggregate line.

`ROLLUP` has a more complex effect when there are multiple `GROUP BY` columns. In this case, each time there is a change in value in any but the last grouping column, the query produces an extra super-aggregate summary row.

For example, without `ROLLUP`, a summary of the `sales` table based on `year`, `country`, and `product` might look like this, where the output indicates summary values only at the year/country/product level of analysis:

```
mysql> SELECT year, country, product, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC, country ASC, product ASC;
+------+---------+------------+--------+
| year | country | product    | profit |
+------+---------+------------+--------+
| 2000 | Finland | Computer   |   1500 |
| 2000 | Finland | Phone      |    100 |
| 2000 | India   | Calculator |    150 |
| 2000 | India   | Computer   |   1200 |
| 2000 | USA     | Calculator |     75 |
| 2000 | USA     | Computer   |   1500 |
| 2001 | Finland | Phone      |     10 |
| 2001 | USA     | Calculator |     50 |
| 2001 | USA     | Computer   |   2700 |
| 2001 | USA     | TV         |    250 |
+------+---------+------------+--------+
```

With `ROLLUP` added, the query produces several extra rows:

```
mysql> SELECT year, country, product, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC, country ASC, product ASC WITH ROLLUP;
+------+---------+------------+--------+
| year | country | product    | profit |
+------+---------+------------+--------+
| 2000 | Finland | Computer   |   1500 |
| 2000 | Finland | Phone      |    100 |
| 2000 | Finland | NULL       |   1600 |
| 2000 | India   | Calculator |    150 |
| 2000 | India   | Computer   |   1200 |
| 2000 | India   | NULL       |   1350 |
| 2000 | USA     | Calculator |     75 |
| 2000 | USA     | Computer   |   1500 |
```

```
| 2000 | USA     | NULL       |   1575 |
| 2000 | NULL    | NULL       |   4525 |
| 2001 | Finland | Phone      |     10 |
| 2001 | Finland | NULL       |     10 |
| 2001 | USA     | Calculator |     50 |
| 2001 | USA     | Computer   |   2700 |
| 2001 | USA     | TV         |    250 |
| 2001 | USA     | NULL       |   3000 |
| 2001 | NULL    | NULL       |   3010 |
| NULL | NULL    | NULL       |   7535 |
+------+---------+------------+--------+
```

Now the output includes summary information at four levels of analysis, not just one:

- Following each set of product rows for a given year and country, an extra summary row appears showing the total for all products. These rows have the `product` column set to `NULL`.
- Following each set of rows for a given year, an extra summary row appears showing the total for all countries and products. These rows have the `country` and `products` columns set to `NULL`.
- Finally, following all other rows, an extra summary row appears showing the grand total for all years, countries, and products. This row has the `year`, `country`, and `products` columns set to `NULL`.

The `GROUPING()` function is available for use in the select list or `HAVING` clause of `GROUP BY ... WITH ROLLUP` queries, to test whether `NULL` values in the result represent super-aggregate values. For example, `GROUPING(year)` returns 1 when `NULL` in the `year` column occurs in a super-aggregate row, and 0 otherwise. Similarly, `GROUPING(country)` and `GROUPING(product)` return 1 for super-aggregate `NULL` values in the `country` and `product` columns, respectively:

```
mysql> SELECT
         year, country, product, SUM(profit) AS profit,
         GROUPING(year) AS grp_year,
         GROUPING(country) AS grp_country,
         GROUPING(product) AS grp_product
       FROM sales
       GROUP BY year ASC, country ASC, product ASC WITH ROLLUP;
+------+---------+------------+--------+----------+-------------+-------------+
| year | country | product    | profit | grp_year | grp_country | grp_product |
+------+---------+------------+--------+----------+-------------+-------------+
| 2000 | Finland | Computer   |   1500 |        0 |           0 |           0 |
| 2000 | Finland | Phone      |    100 |        0 |           0 |           0 |
| 2000 | Finland | NULL       |   1600 |        0 |           0 |           1 |
| 2000 | India   | Calculator |    150 |        0 |           0 |           0 |
| 2000 | India   | Computer   |   1200 |        0 |           0 |           0 |
| 2000 | India   | NULL       |   1350 |        0 |           0 |           1 |
| 2000 | USA     | Calculator |     75 |        0 |           0 |           0 |
| 2000 | USA     | Computer   |   1500 |        0 |           0 |           0 |
| 2000 | USA     | NULL       |   1575 |        0 |           0 |           1 |
| 2000 | NULL    | NULL       |   4525 |        0 |           1 |           1 |
| 2001 | Finland | Phone      |     10 |        0 |           0 |           0 |
| 2001 | Finland | NULL       |     10 |        0 |           0 |           1 |
| 2001 | USA     | Calculator |     50 |        0 |           0 |           0 |
| 2001 | USA     | Computer   |   2700 |        0 |           0 |           0 |
| 2001 | USA     | TV         |    250 |        0 |           0 |           0 |
| 2001 | USA     | NULL       |   3000 |        0 |           0 |           1 |
| 2001 | NULL    | NULL       |   3010 |        0 |           1 |           1 |
| NULL | NULL    | NULL       |   7535 |        1 |           1 |           1 |
+------+---------+------------+--------+----------+-------------+-------------+
```

Instead of displaying the `GROUPING()` results directly, you can use `GROUPING()` to substitute labels for super-aggregate `NULL` values:

```
mysql> SELECT
         IF(GROUPING(year), 'All years', year) AS year,
         IF(GROUPING(country), 'All countries', country) AS country,
         IF(GROUPING(product), 'All products', product) AS product,
         SUM(profit) AS profit
```

```
       FROM sales
       GROUP BY year ASC, country ASC, product ASC WITH ROLLUP;
+-----------+---------------+-------------+--------+
| year      | country       | product     | profit |
+-----------+---------------+-------------+--------+
| 2000      | Finland       | Computer    |   1500 |
| 2000      | Finland       | Phone       |    100 |
| 2000      | Finland       | All products |   1600 |
| 2000      | India         | Calculator  |    150 |
| 2000      | India         | Computer    |   1200 |
| 2000      | India         | All products |   1350 |
| 2000      | USA           | Calculator  |     75 |
| 2000      | USA           | Computer    |   1500 |
| 2000      | USA           | All products |   1575 |
| 2000      | All countries | All products |   4525 |
| 2001      | Finland       | Phone       |     10 |
| 2001      | Finland       | All products |     10 |
| 2001      | USA           | Calculator  |     50 |
| 2001      | USA           | Computer    |   2700 |
| 2001      | USA           | TV          |    250 |
| 2001      | USA           | All products |   3000 |
| 2001      | All countries | All products |   3010 |
| All years | All countries | All products |   7535 |
+-----------+---------------+-------------+--------+
```

With multiple expression arguments, `GROUPING()` returns a result representing a bitmask the combines the results for each expression, with the lowest-order bit corresponding to the result for the rightmost expression. For example, `GROUPING(year, country, product)` is evaluated like this:

```
  result for GROUPING(product)
+ result for GROUPING(country) << 1
+ result for GROUPING(year) << 2
```

The result of such a `GROUPING()` is nonzero if any of the expressions represents a super-aggregate `NULL`, so you can return only the super-aggregate rows and filter out the regular grouped rows like this:

```
mysql> SELECT year, country, product, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC, country ASC, product ASC WITH ROLLUP
       HAVING GROUPING(year, country, product) <> 0;
+------+---------+---------+--------+
| year | country | product | profit |
+------+---------+---------+--------+
| 2000 | Finland | NULL    |   1600 |
| 2000 | India   | NULL    |   1350 |
| 2000 | USA     | NULL    |   1575 |
| 2000 | NULL    | NULL    |   4525 |
| 2001 | Finland | NULL    |     10 |
| 2001 | USA     | NULL    |   3000 |
| 2001 | NULL    | NULL    |   3010 |
| NULL | NULL    | NULL    |   7535 |
+------+---------+---------+--------+
```

The `sales` table contains no `NULL` values, so all `NULL` values in a `ROLLUP` result represent super-aggregate values. When the data set contains `NULL` values, `ROLLUP` summaries may contain `NULL` values not only in super-aggregate rows, but also in regular grouped rows. `GROUPING()` enables these to be distinguished. Suppose that table `t1` contains a simple data set with two grouping factors for a set of quantity values, where `NULL` indicates something like “other” or “unknown”:

```
mysql> SELECT * FROM t1;
+------+-------+----------+
| name | size  | quantity |
+------+-------+----------+
| ball | small |       10 |
| ball | large |       20 |
| ball | NULL  |        5 |
| hoop | small |       15 |
```

```
| hoop | large |        5 |
| hoop | NULL  |        3 |
+------+-------+----------+
```

A simple `ROLLUP` operation produces these results, in which it is not so easy to distinguish `NULL` values in super-aggregate rows from `NULL` values in regular grouped rows:

```
mysql> SELECT name, size, SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+------+-------+----------+
| name | size  | quantity |
+------+-------+----------+
| ball | NULL  |        5 |
| ball | large |       20 |
| ball | small |       10 |
| ball | NULL  |       35 |
| hoop | NULL  |        3 |
| hoop | large |        5 |
| hoop | small |       15 |
| hoop | NULL  |       23 |
| NULL | NULL  |       58 |
+------+-------+----------+
```

Using `GROUPING()` to substitute labels for the super-aggregate `NULL` values makes the result easier to interpret:

```
mysql> SELECT
         IF(GROUPING(name) = 1, 'All items', name) AS name,
         IF(GROUPING(size) = 1, 'All sizes', size) AS size,
         SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+-----------+-----------+----------+
| name      | size      | quantity |
+-----------+-----------+----------+
| ball      | NULL      |        5 |
| ball      | large     |       20 |
| ball      | small     |       10 |
| ball      | All sizes |       35 |
| hoop      | NULL      |        3 |
| hoop      | large     |        5 |
| hoop      | small     |       15 |
| hoop      | All sizes |       23 |
| All items | All sizes |       58 |
+-----------+-----------+----------+
```

Other Considerations When using ROLLUP

The following discussion lists some behaviors specific to the MySQL implementation of `ROLLUP`.

When you use `ROLLUP`, you cannot also use an `ORDER BY` clause to sort the results. In other words, `ROLLUP` and `ORDER BY` are mutually exclusive in MySQL. However, you still have some control over sort order. `GROUP BY` in MySQL implicitly sorts results by default (in the absence of `ASC` or `DESC` designators). However, implicit `GROUP BY` sorting in MySQL is deprecated. To achieve a specific sort order of grouped results:

- Use explicit `ASC` and `DESC` keywords with columns named in the `GROUP BY` list to specify sort order for individual columns. In this case, the super-aggregate summary rows added by `ROLLUP` still appear after the rows from which they are calculated, regardless of the sort order.
- To work around the restriction that prevents using `ROLLUP` with `ORDER BY`, generate the grouped result set as a derived table and apply `ORDER BY` to it. For example:

```
SELECT * FROM
```

```
(SELECT year, SUM(profit) FROM sales GROUP BY year WITH ROLLUP) AS dt
ORDER BY year;
```

In this case, the super-aggregate summary rows sort with the rows from which they are calculated, and their placement depends on sort order (at the beginning for ascending sort, at the end for descending sort).

`LIMIT` can be used to restrict the number of rows returned to the client. `LIMIT` is applied after `ROLLUP`, so the limit applies against the extra rows added by `ROLLUP`. For example:

```
mysql> SELECT year, country, product, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC, country ASC, product ASC WITH ROLLUP
       LIMIT 5;
+------+---------+------------+--------+
| year | country | product    | profit |
+------+---------+------------+--------+
| 2000 | Finland | Computer   |   1500 |
| 2000 | Finland | Phone      |    100 |
| 2000 | Finland | NULL       |   1600 |
| 2000 | India   | Calculator |    150 |
| 2000 | India   | Computer   |   1200 |
+------+---------+------------+--------+
```

Using `LIMIT` with `ROLLUP` may produce results that are more difficult to interpret, because there is less context for understanding the super-aggregate rows.

The `NULL` indicators in each super-aggregate row are produced when the row is sent to the client. The server looks at the columns named in the `GROUP BY` clause following the leftmost one that has changed value. For any column in the result set with a name that matches any of those names, its value is set to `NULL`. (If you specify grouping columns by column position, the server identifies which columns to set to `NULL` by position.)

Because the `NULL` values in the super-aggregate rows are placed into the result set at such a late stage in query processing, you can test them as `NULL` values only in the select list or `HAVING` clause. You cannot test them as `NULL` values in join conditions or the `WHERE` clause to determine which rows to select. For example, you cannot add `WHERE product IS NULL` to the query to eliminate from the output all but the super-aggregate rows.

The `NULL` values do appear as `NULL` on the client side and can be tested as such using any MySQL client programming interface. However, at this point, you cannot distinguish whether a `NULL` represents a regular grouped value or a super-aggregate value.

A MySQL extension permits a column that does not appear in the `GROUP BY` list to be named in the select list. (For information about nonaggregated columns and `GROUP BY`, see Section 12.18.3, “MySQL Handling of GROUP BY”.) In this case, the server is free to choose any value from this nonaggregated column in summary rows, and this includes the extra rows added by `WITH ROLLUP`. For example, in the following query, `country` is a nonaggregated column that does not appear in the `GROUP BY` list and values chosen for this column are indeterminate:

```
mysql> SELECT year, country, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC WITH ROLLUP;
+------+---------+--------+
| year | country | profit |
+------+---------+--------+
| 2000 | India   |   4525 |
| 2001 | USA     |   3010 |
| NULL | USA     |   7535 |
+------+---------+--------+
```

This behavior is permitted when the `ONLY_FULL_GROUP_BY` SQL mode is not enabled. If that mode is enabled, the server rejects the query as illegal because `country` is not listed in the `GROUP BY` clause.

With ONLY_FULL_GROUP_BY enabled, you can still execute the query by using the ANY_VALUE() function for indeterminate-value columns:

```
mysql> SELECT year, ANY_VALUE(country) AS country, SUM(profit) AS profit
       FROM sales
       GROUP BY year ASC WITH ROLLUP;
+------+---------+--------+
| year | country | profit |
+------+---------+--------+
| 2000 | India   |   4525 |
| 2001 | USA     |   3010 |
| NULL | USA     |   7535 |
+------+---------+--------+
```

12.18.3 MySQL Handling of GROUP BY

SQL92 and earlier does not permit queries for which the select list, HAVING condition, or ORDER BY list refer to nonaggregated columns that are not named in the GROUP BY clause. For example, this query is illegal in standard SQL92 because the nonaggregated name column in the select list does not appear in the GROUP BY:

```
SELECT o.custid, c.name, MAX(o.payment)
  FROM orders AS o, customers AS c
  WHERE o.custid = c.custid
  GROUP BY o.custid;
```

For the query to be legal in SQL92, the name column must be omitted from the select list or named in the GROUP BY clause.

SQL99 and later permits such nonaggregates per optional feature T301 if they are functionally dependent on GROUP BY columns: If such a relationship exists between name and custid, the query is legal. This would be the case, for example, were custid a primary key of customers.

MySQL implements detection of functional dependence. If the ONLY_FULL_GROUP_BY SQL mode is enabled (which it is by default), MySQL rejects queries for which the select list, HAVING condition, or ORDER BY list refer to nonaggregated columns that are neither named in the GROUP BY clause nor are functionally dependent on them.

If ONLY_FULL_GROUP_BY is disabled, a MySQL extension to the standard SQL use of GROUP BY permits the select list, HAVING condition, or ORDER BY list to refer to nonaggregated columns even if the columns are not functionally dependent on GROUP BY columns. This causes MySQL to accept the preceding query. In this case, the server is free to choose any value from each group, so unless they are the same, the values chosen are indeterminate, which is probably not what you want. Furthermore, the selection of values from each group cannot be influenced by adding an ORDER BY clause. Result set sorting occurs after values have been chosen, and ORDER BY does not affect which value within each group the server chooses. Disabling ONLY_FULL_GROUP_BY is useful primarily when you know that, due to some property of the data, all values in each nonaggregated column not named in the GROUP BY are the same for each group.

You can achieve the same effect without disabling ONLY_FULL_GROUP_BY by using ANY_VALUE() to refer to the nonaggregated column.

The following discussion demonstrates functional dependence, the error message MySQL produces when functional dependence is absent, and ways of causing MySQL to accept a query in the absence of functional dependence.

This query might be invalid with ONLY_FULL_GROUP_BY enabled because the nonaggregated address column in the select list is not named in the GROUP BY clause:

```
SELECT name, address, MAX(age) FROM t GROUP BY name;
```

The query is valid if `name` is a primary key of `t` or is a unique `NOT NULL` column. In such cases, MySQL recognizes that the selected column is functionally dependent on a grouping column. For example, if `name` is a primary key, its value determines the value of `address` because each group has only one value of the primary key and thus only one row. As a result, there is no randomness in the choice of `address` value in a group and no need to reject the query.

The query is invalid if `name` is not a primary key of `t` or a unique `NOT NULL` column. In this case, no functional dependency can be inferred and an error occurs:

```
mysql> SELECT name, address, MAX(age) FROM t GROUP BY name;
ERROR 1055 (42000): Expression #2 of SELECT list is not in GROUP
BY clause and contains nonaggregated column 'mydb.t.address' which
is not functionally dependent on columns in GROUP BY clause; this
is incompatible with sql_mode=only_full_group_by
```

If you know that, *for a given data set,* each `name` value in fact uniquely determines the `address` value, `address` is effectively functionally dependent on `name`. To tell MySQL to accept the query, you can use the `ANY_VALUE()` function:

```
SELECT name, ANY_VALUE(address), MAX(age) FROM t GROUP BY name;
```

Alternatively, disable `ONLY_FULL_GROUP_BY`.

The preceding example is quite simple, however. In particular, it is unlikely you would group on a single primary key column because every group would contain only one row. For addtional examples demonstrating functional dependence in more complex queries, see Section 12.18.4, “Detection of Functional Dependence”.

If a query has aggregate functions and no `GROUP BY` clause, it cannot have nonaggregated columns in the select list, `HAVING` condition, or `ORDER BY` list with `ONLY_FULL_GROUP_BY` enabled:

```
mysql> SELECT name, MAX(age) FROM t;
ERROR 1140 (42000): In aggregated query without GROUP BY, expression
#1 of SELECT list contains nonaggregated column 'mydb.t.name'; this
is incompatible with sql_mode=only_full_group_by
```

Without `GROUP BY`, there is a single group and it is indeterminate which `name` value to choose for the group. Here, too, `ANY_VALUE()` can be used, if it is immaterial which `name` value MySQL chooses:

```
SELECT ANY_VALUE(name), MAX(age) FROM t;
```

`ONLY_FULL_GROUP_BY` also affects handling of queries that use `DISTINCT` and `ORDER BY`. Consider the case of a table `t` with three columns `c1`, `c2`, and `c3` that contains these rows:

```
c1 c2 c3
1  2  A
3  4  B
1  2  C
```

Suppose that we execute the following query, expecting the results to be ordered by `c3`:

```
SELECT DISTINCT c1, c2 FROM t ORDER BY c3;
```

To order the result, duplicates must be eliminated first. But to do so, should we keep the first row or the third? This arbitrary choice influences the retained value of `c3`, which in turn influences ordering

and makes it arbitrary as well. To prevent this problem, a query that has `DISTINCT` and `ORDER BY` is rejected as invalid if any `ORDER BY` expression does not satisfy at least one of these conditions:

- The expression is equal to one in the select list
- All columns referenced by the expression and belonging to the query's selected tables are elements of the select list

Another MySQL extension to standard SQL permits references in the `HAVING` clause to aliased expressions in the select list. For example, the following query returns `name` values that occur only once in table `orders`:

```
SELECT name, COUNT(name) FROM orders
  GROUP BY name
  HAVING COUNT(name) = 1;
```

The MySQL extension permits the use of an alias in the `HAVING` clause for the aggregated column:

```
SELECT name, COUNT(name) AS c FROM orders
  GROUP BY name
  HAVING c = 1;
```

Standard SQL permits only column expressions in `GROUP BY` clauses, so a statement such as this is invalid because `FLOOR(value/100)` is a noncolumn expression:

```
SELECT id, FLOOR(value/100)
  FROM tbl_name
  GROUP BY id, FLOOR(value/100);
```

MySQL extends standard SQL to permit noncolumn expressions in `GROUP BY` clauses and considers the preceding statement valid.

Standard SQL also does not permit aliases in `GROUP BY` clauses. MySQL extends standard SQL to permit aliases, so another way to write the query is as follows:

```
SELECT id, FLOOR(value/100) AS val
  FROM tbl_name
  GROUP BY id, val;
```

The alias `val` is considered a column expression in the `GROUP BY` clause.

In the presence of a noncolumn expression in the `GROUP BY` clause, MySQL recognizes equality between that expression and expressions in the select list. This means that with `ONLY_FULL_GROUP_BY` SQL mode enabled, the query containing `GROUP BY id, FLOOR(value/100)` is valid because that same `FLOOR()` expression occurs in the select list. However, MySQL does not try to recognize functional dependence on `GROUP BY` noncolumn expressions, so the following query is invalid with `ONLY_FULL_GROUP_BY` enabled, even though the third selected expression is a simple formula of the `id` column and the `FLOOR()` expression in the `GROUP BY` clause:

```
SELECT id, FLOOR(value/100), id+FLOOR(value/100)
  FROM tbl_name
  GROUP BY id, FLOOR(value/100);
```

A workaround is to use a derived table:

```
SELECT id, F, id+F
  FROM
    (SELECT id, FLOOR(value/100) AS F
```

```
    FROM tbl_name
    GROUP BY id, FLOOR(value/100)) AS dt;
```

12.18.4 Detection of Functional Dependence

The following discussion provides several examples of the ways in which MySQL detects functional dependencies. The examples use this notation:

```
{X} -> {Y}
```

Understand this as "*X* uniquely determines *Y*," which also means that *Y* is functionally dependent on *X*.

The examples use the `world` database, which can be downloaded from the MySQL Documentation page. You can find details on how to install the database on the same page.

- Functional Dependencies Derived from Keys
- Functional Dependencies Derived from Multiple-Column Keys and from Equalities
- Functional Dependency Special Cases
- Functional Dependencies and Views
- Combinations of Functional Dependencies

Functional Dependencies Derived from Keys

The following query selects, for each country, a count of spoken languages:

```
SELECT co.Name, COUNT(*)
FROM countrylanguage cl, country co
WHERE cl.CountryCode = co.Code
GROUP BY co.Code;
```

`co.Code` is a primary key of `co`, so all columns of `co` are functionally dependent on it, as expressed using this notation:

```
{co.Code} -> {co.*}
```

Thus, `co.name` is functionally dependent on `GROUP BY` columns and the query is valid.

A `UNIQUE` index over a `NOT NULL` column could be used instead of a primary key and the same functional dependence would apply. (This is not true for a `UNIQUE` index that permits `NULL` values because it permits multiple `NULL` values and in that case uniqueness is lost.)

Functional Dependencies Derived from Multiple-Column Keys and from Equalities

This query selects, for each country, a list of all spoken languages and how many people speak them:

```
SELECT co.Name, cl.Language,
cl.Percentage * co.Population / 100.0 AS SpokenBy
FROM countrylanguage cl, country co
WHERE cl.CountryCode = co.Code
GROUP BY cl.CountryCode, cl.Language;
```

The pair (`cl.CountryCode`, `cl.Language`) is a two-column composite primary key of `cl`, so that column pair uniquely determines all columns of `cl`:

```
{cl.CountryCode, cl.Language} -> {cl.*}
```

Moreover, because of the equality in the `WHERE` clause:

```
{cl.CountryCode} -> {co.Code}
```

And, because `co.Code` is primary key of `co`:

```
{co.Code} -> {co.*}
```

“Uniquely determines” relationships are transitive, therefore:

```
{cl.CountryCode, cl.Language} -> {cl.*,co.*}
```

As a result, the query is valid.

As with the previous example, a `UNIQUE` key over `NOT NULL` columns could be used instead of a primary key.

An `INNER JOIN` condition can be used instead of `WHERE`. The same functional dependencies apply:

```
SELECT co.Name, cl.Language,
cl.Percentage * co.Population/100.0 AS SpokenBy
FROM countrylanguage cl INNER JOIN country co
ON cl.CountryCode = co.Code
GROUP BY cl.CountryCode, cl.Language;
```

Functional Dependency Special Cases

Whereas an equality test in a `WHERE` condition or `INNER JOIN` condition is symmetric, an equality test in an outer join condition is not, because tables play different roles.

Assume that referential integrity has been accidentally broken and there exists a row of `countrylanguage` without a corresponding row in `country`. Consider the same query as in the previous example, but with a `LEFT JOIN`:

```
SELECT co.Name, cl.Language,
cl.Percentage * co.Population/100.0 AS SpokenBy
FROM countrylanguage cl LEFT JOIN country co
ON cl.CountryCode = co.Code
GROUP BY cl.CountryCode, cl.Language;
```

For a given value of `cl.CountryCode`, the value of `co.Code` in the join result is either found in a matching row (determined by `cl.CountryCode`) or is `NULL`-complemented if there is no match (also determined by `cl.CountryCode`). In each case, this relationship applies:

```
{cl.CountryCode} -> {co.Code}
```

`cl.CountryCode` is itself functionally dependent on {`cl.CountryCode`, `cl.Language`} which is a primary key.

If in the join result `co.Code` is `NULL`-complemented, `co.Name` is as well. If `co.Code` is not `NULL`-complemented, then because `co.Code` is a primary key, it determines `co.Name`. Therefore, in all cases:

```
{co.Code} -> {co.Name}
```

Which yields:

```
{cl.CountryCode, cl.Language} -> {cl.*,co.*}
```

As a result, the query is valid.

However, suppose that the tables are swapped, as in this query:

```
SELECT co.Name, cl.Language,
cl.Percentage * co.Population/100.0 AS SpokenBy
FROM country co LEFT JOIN countrylanguage cl
ON cl.CountryCode = co.Code
GROUP BY cl.CountryCode, cl.Language;
```

Now this relationship does *not* apply:

```
{cl.CountryCode, cl.Language} -> {cl.*,co.*}
```

Indeed, all `NULL`-complemented rows made for `cl` will be put into a single group (they have both `GROUP BY` columns equal to `NULL`), and inside this group the value of `co.Name` can vary. The query is invalid and MySQL rejects it.

Functional dependence in outer joins is thus linked to whether determinant columns belong to the left or right side of the `LEFT JOIN`. Determination of functional dependence becomes more complex if there are nested outer joins or the join condition does not consist entirely of equality comparisons.

Functional Dependencies and Views

Suppose that a view on countries produces their code, their name in uppercase, and how many different official languages they have:

```
CREATE VIEW Country2 AS
SELECT co.Code, UPPER(co.Name) AS UpperName,
COUNT(cl.Language) AS OfficialLanguages
FROM country AS co JOIN countrylanguage AS cl
ON cl.CountryCode = co.Code
WHERE cl.isOfficial = 'T'
GROUP BY co.Code;
```

This definition is valid because:

```
{co.Code} -> {co.*}
```

In the view result, the first selected column is `co.Code`, which is also the group column and thus determines all other selected expressions:

```
{Country2.Code} -> {Country2.*}
```

MySQL understands this and uses this information, as described following.

This query displays countries, how many different official languages they have, and how many cities they have, by joining the view with the `city` table:

```
SELECT co2.Code, co2.UpperName, co2.OfficialLanguages,
COUNT(*) AS Cities
FROM country2 AS co2 JOIN city ci
ON ci.CountryCode = co2.Code
GROUP BY co2.Code;
```

This query is valid because, as seen previously:

```
```

```
{co2.Code} -> {co2.*}
```

MySQL is able to discover a functional dependency in the result of a view and use that to validate a query which uses the view. The same would be true if `country2` were a derived table (or common table expression), as in:

```
SELECT co2.Code, co2.UpperName, co2.OfficialLanguages,
COUNT(*) AS Cities
FROM
(
 SELECT co.Code, UPPER(co.Name) AS UpperName,
 COUNT(cl.Language) AS OfficialLanguages
 FROM country AS co JOIN countrylanguage AS cl
 ON cl.CountryCode=co.Code
 WHERE cl.isOfficial='T'
 GROUP BY co.Code
) AS co2
JOIN city ci ON ci.CountryCode = co2.Code
GROUP BY co2.Code;
```

Combinations of Functional Dependencies

MySQL is able to combine all of the preceding types of functional dependencies (key based, equality based, view based) to validate more complex queries.

12.19 Internal Functions

Table 12.24 Internal Functions

Name	Description
`CAN_ACCESS_COLUMN()`	Internal use only
`CAN_ACCESS_DATABASE()`	Internal use only
`CAN_ACCESS_TABLE()`	Internal use only
`CAN_ACCESS_VIEW()`	Internal use only
`GET_DD_COLUMN_PRIVILEGES()`	Internal use only
`GET_DD_CREATE_OPTIONS()`	Internal use only
`GET_DD_INDEX_SUB_PART_LENGTH()`	Internal use only
`INTERNAL_AUTO_INCREMENT()`	Internal use only
`INTERNAL_AVG_ROW_LENGTH()`	Internal use only
`INTERNAL_CHECK_TIME()`	Internal use only
`INTERNAL_CHECKSUM()`	Internal use only
`INTERNAL_DATA_FREE()`	Internal use only
`INTERNAL_DATA_LENGTH()`	Internal use only
`INTERNAL_DD_CHAR_LENGTH()`	Internal use only
`INTERNAL_GET_COMMENT_OR_ERROR()`	Internal use only
`INTERNAL_GET_VIEW_WARNING_OR_ERROR()`	Internal use only
`INTERNAL_INDEX_COLUMN_CARDINALITY()`	Internal use only
`INTERNAL_INDEX_LENGTH()`	Internal use only
`INTERNAL_KEYS_DISABLED()`	Internal use only
`INTERNAL_MAX_DATA_LENGTH()`	Internal use only
`INTERNAL_TABLE_ROWS()`	Internal use only
`INTERNAL_UPDATE_TIME()`	Internal use only

The functions listed in this section are intended only for internal use by the server. Attempts by users to invoke them result in an error.

- `CAN_ACCESS_COLUMN(ARGS)`
- `CAN_ACCESS_DATABASE(ARGS)`
- `CAN_ACCESS_TABLE(ARGS)`
- `CAN_ACCESS_VIEW(ARGS)`
- `GET_DD_COLUMN_PRIVILEGES(ARGS)`
- `GET_DD_CREATE_OPTIONS(ARGS)`
- `GET_DD_INDEX_SUB_PART_LENGTH(ARGS)`
- `INTERNAL_AUTO_INCREMENT(ARGS)`
- `INTERNAL_AVG_ROW_LENGTH(ARGS)`
- `INTERNAL_CHECK_TIME(ARGS)`
- `INTERNAL_CHECKSUM(ARGS)`
- `INTERNAL_DATA_FREE(ARGS)`
- `INTERNAL_DATA_LENGTH(ARGS)`
- `INTERNAL_DD_CHAR_LENGTH(ARGS)`
- `INTERNAL_GET_COMMENT_OR_ERROR(ARGS)`
- `INTERNAL_GET_VIEW_WARNING_OR_ERROR(ARGS)`
- `INTERNAL_INDEX_COLUMN_CARDINALITY(ARGS)`
- `INTERNAL_INDEX_LENGTH(ARGS)`
- `INTERNAL_KEYS_DISABLED(ARGS)`
- `INTERNAL_MAX_DATA_LENGTH(ARGS)`
- `INTERNAL_TABLE_ROWS(ARGS)`
- `INTERNAL_UPDATE_TIME(ARGS)`
- `IS_VISIBLE_DD_OBJECT(ARGS)`

12.20 Miscellaneous Functions

Table 12.25 Miscellaneous Functions

Name	Description
`ANY_VALUE()`	Suppress ONLY_FULL_GROUP_BY value rejection
`BIN_TO_UUID()`	Convert binary UUID to string
`DEFAULT()`	Return the default value for a table column
`GET_LOCK()`	Get a named lock
`GROUPING()`	Distinguish super-aggregate ROLLUP rows from regular rows
`INET_ATON()`	Return the numeric value of an IP address

Name	Description
INET_NTOA()	Return the IP address from a numeric value
INET6_ATON()	Return the numeric value of an IPv6 address
INET6_NTOA()	Return the IPv6 address from a numeric value
IS_FREE_LOCK()	Whether the named lock is free
IS_IPV4()	Whether argument is an IPv4 address
IS_IPV4_COMPAT()	Whether argument is an IPv4-compatible address
IS_IPV4_MAPPED()	Whether argument is an IPv4-mapped address
IS_IPV6()	Whether argument is an IPv6 address
IS_USED_LOCK()	Whether the named lock is in use; return connection identifier if true
IS_UUID()	Whether argument is a valid UUID
MASTER_POS_WAIT()	Block until the slave has read and applied all updates up to the specified position
NAME_CONST()	Causes the column to have the given name
RAND()	Return a random floating-point value
RELEASE_ALL_LOCKS()	Releases all current named locks
RELEASE_LOCK()	Releases the named lock
SLEEP()	Sleep for a number of seconds
UUID()	Return a Universal Unique Identifier (UUID)
UUID_SHORT()	Return an integer-valued universal identifier
UUID_TO_BIN()	Convert string UUID to binary
VALUES()	Defines the values to be used during an INSERT

- ANY_VALUE(*arg*)

  This function is useful for GROUP BY queries when the ONLY_FULL_GROUP_BY SQL mode is enabled, for cases when MySQL rejects a query that you know is valid for reasons that MySQL cannot determine. The function return value and type are the same as the return value and type of its argument, but the function result is not checked for the ONLY_FULL_GROUP_BY SQL mode.

  For example, if name is a nonindexed column, the following query fails with ONLY_FULL_GROUP_BY enabled:

```
mysql> SELECT name, address, MAX(age) FROM t GROUP BY name;
ERROR 1055 (42000): Expression #2 of SELECT list is not in GROUP
BY clause and contains nonaggregated column 'mydb.t.address' which
is not functionally dependent on columns in GROUP BY clause; this
is incompatible with sql_mode=only_full_group_by
```

  The failure occurs because address is a nonaggregated column that is neither named among GROUP BY columns nor functionally dependent on them. As a result, the address value for rows within each name group is nondeterministic. There are multiple ways to cause MySQL to accept the query:

  - Alter the table to make name a primary key or a unique NOT NULL column. This enables MySQL to determine that address is functionally dependent on name; that is, address is uniquely determined by name. (This technique is inapplicable if NULL must be permitted as a valid name value.)
  - Use ANY_VALUE() to refer to address:

```
SELECT name, ANY_VALUE(address), MAX(age) FROM t GROUP BY name;
```

In this case, MySQL ignores the nondeterminism of `address` values within each `name` group and accepts the query. This may be useful if you simply do not care which value of a nonaggregated column is chosen for each group. `ANY_VALUE()` is not an aggregate function, unlike functions such as `SUM()` or `COUNT()`. It simply acts to suppress the test for nondeterminism.

- Disable `ONLY_FULL_GROUP_BY`. This is equivalent to using `ANY_VALUE()` with `ONLY_FULL_GROUP_BY` enabled, as described in the previous item.

`ANY_VALUE()` is also useful if functional dependence exists between columns but MySQL cannot determine it. The following query is valid because `age` is functionally dependent on the grouping column `age-1`, but MySQL cannot tell that and rejects the query with `ONLY_FULL_GROUP_BY` enabled:

```
SELECT age FROM t GROUP BY age-1;
```

To cause MySQL to accept the query, use `ANY_VALUE()`:

```
SELECT ANY_VALUE(age) FROM t GROUP BY age-1;
```

`ANY_VALUE()` can be used for queries that refer to aggregate functions in the absence of a `GROUP BY` clause:

```
mysql> SELECT name, MAX(age) FROM t;
ERROR 1140 (42000): In aggregated query without GROUP BY, expression
#1 of SELECT list contains nonaggregated column 'mydb.t.name'; this
is incompatible with sql_mode=only_full_group_by
```

Without `GROUP BY`, there is a single group and it is indeterminate which `name` value to choose for the group. `ANY_VALUE()` tells MySQL to accept the query:

```
SELECT ANY_VALUE(name), MAX(age) FROM t;
```

It may be that, due to some property of a given data set, you know that a selected nonaggregated column is effectively functionally dependent on a `GROUP BY` column. For example, an application may enforce uniqueness of one column with respect to another. In this case, using `ANY_VALUE()` for the effectively functionally dependent column may make sense.

For additional discussion, see Section 12.18.3, “MySQL Handling of GROUP BY”.

- `BIN_TO_UUID(binary_uuid)`, `BIN_TO_UUID(binary_uuid, swap_flag)`

  `BIN_TO_UUID()` is the inverse of `UUID_TO_BIN()`. It converts a binary UUID to a string UUID and returns the result. The binary value should be a UUID as a `VARBINARY(16)` value. The return value is a `utf8` string of five hexadecimal numbers separated by dashes. (For details about this format, see the `UUID()` function description.) If the UUID argument is `NULL`, the return value is `NULL`. If any argument is invalid, an error occurs.

  `BIN_TO_UUID()` takes one or two arguments:

  - The one-argument form takes a binary UUID value. The UUID value is assumed not to have its time-low and time-high parts swapped. The string result is in the same order as the binary argument.
  - The two-argument form takes a binary UUID value and a swap-flag value:

- If *swap_flag* is 0, the two-argument form is equivalent to the one-argument form. The string result is in the same order as the binary argument.
- If *swap_flag* is 1, the UUID value is assumed to have its time-low and time-high parts swapped. These parts are swapped back to their original position in the result value.

For usage examples and information about time-part swapping, see the `UUID_TO_BIN()` function description.

- `DEFAULT(col_name)`

  Returns the default value for a table column. An error results if the column has no default value.

```
mysql> UPDATE t SET i = DEFAULT(i)+1 WHERE id < 100;
```

- `FORMAT(X,D)`

  Formats the number *X* to a format like `'#,###,###.##'`, rounded to *D* decimal places, and returns the result as a string. For details, see Section 12.5, “String Functions”.

- `GET_LOCK(str,timeout)`

  Tries to obtain a lock with a name given by the string *str*, using a timeout of *timeout* seconds. A negative *timeout* value means infinite timeout. The lock is exclusive. While held by one session, other sessions cannot obtain a lock of the same name.

  Returns `1` if the lock was obtained successfully, `0` if the attempt timed out (for example, because another client has previously locked the name), or `NULL` if an error occurred (such as running out of memory or the thread was killed with `mysqladmin kill`).

  A lock obtained with `GET_LOCK()` is released explicitly by executing `RELEASE_LOCK()` or implicitly when your session terminates (either normally or abnormally). Lock release may also occur with another call to `GET_LOCK()`:

  - `GET_LOCK()` is implemented using the metadata locking (MDL) subsystem. Multiple simultaneous locks can be acquired and `GET_LOCK()` does not release any existing locks. It is even possible for a given session to acquire multiple locks for the same name. Other sessions cannot acquire a lock with that name until the acquiring session releases all its locks for the name.

    Locks acquired with `GET_LOCK()` appear in the Performance Schema `metadata_locks` table. The `OBJECT_TYPE` column says `USER LEVEL LOCK` and the `OBJECT_NAME` column indicates the lock name. Also, the capability of acquiring multiple locks introduces the possibility of deadlock among clients. When this happens, the server chooses a caller and terminates its lock-acquisition request with an `ER_USER_LOCK_DEADLOCK` error. This error does not cause transactions to roll back.

  For example, suppose that you execute these statements:

```
SELECT GET_LOCK('lock1',10);
SELECT GET_LOCK('lock2',10);
SELECT RELEASE_LOCK('lock2');
SELECT RELEASE_LOCK('lock1');
```

  The second `GET_LOCK()` acquires a second lock and both `RELEASE_LOCK()` calls return 1 (success).

  MySQL enforces a maximum length on lock names of 64 characters.

  Locks obtained with `GET_LOCK()` are not released when transactions commit or roll back.

`GET_LOCK()` can be used to implement application locks or to simulate record locks. Names are locked on a server-wide basis. If a name has been locked within one session, `GET_LOCK()` blocks any request by another session for a lock with the same name. This enables clients that agree on a given lock name to use the name to perform cooperative advisory locking. But be aware that it also enables a client that is not among the set of cooperating clients to lock a name, either inadvertently or deliberately, and thus prevent any of the cooperating clients from locking that name. One way to reduce the likelihood of this is to use lock names that are database-specific or application-specific. For example, use lock names of the form *`db_name.str`* or *`app_name.str`*.

If multiple clients are waiting for a lock, the order in which they will acquire it is undefined. Applications should not assume that clients will acquire the lock in the same order that they issued the lock requests.

`GET_LOCK()` is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

Caution

With the capability of acquiring multiple named locks in MySQL 5.7.5, it is possible for a single statement to acquire a large number of locks. For example:

```
INSERT INTO ... SELECT GET_LOCK(t1.col_name) FROM t1;
```

These types of statements may have certain adverse effects. For example, if the statement fails part way through and rolls back, locks acquired up to the point of failure will still exist. If the intent is for there to be a correspondence between rows inserted and locks acquired, that intent will not be satisfied. Also, if it is important that locks are granted in a certain order, be aware that result set order may differ depending on which execution plan the optimizer chooses. For these reasons, it may be best to limit applications to a single lock-acquisition call per statement.

A different locking interface is available as either a plugin service or a set of user-defined functions. This interface provides lock namespaces and distinct read and write locks, unlike the interface provided by `GET_LOCK()` and related functions. For details, see Section 28.3.1, “The Locking Service”.

- `GROUPING(`*`expr [, expr] ...`*`)`

  For `GROUP BY` queries that include a `WITH ROLLUP` modifier, the `ROLLUP` operation produces super-aggregate output rows where `NULL` represents the set of all values. The `GROUPING()` function enables you to distinguish `NULL` values for super-aggregate rows from `NULL` values in regular grouped rows.

  `GROUPING()` is permitted only in the select list or `HAVING` clause.

  Each argument to `GROUPING()` must be an expression that exactly matches an expression in the `GROUP BY` clause. The expression cannot be a positional specifier. For each expression, `GROUPING()` produces 1 if the expression value in the current row is a `NULL` representing a super-aggregate value. Otherwise, `GROUPING()` produces 0, indicating that the expression value is a `NULL` for a regular result row or is not `NULL`.

  Suppose that table `t1` contains these rows, where `NULL` indicates something like “other” or “unknown”:

  ```
  mysql> SELECT * FROM t1;
  +------+-------+----------+
  ```

```
| name | size  | quantity |
+------+-------+----------+
| ball | small |       10 |
| ball | large |       20 |
| ball | NULL  |        5 |
| hoop | small |       15 |
| hoop | large |        5 |
| hoop | NULL  |        3 |
+------+-------+----------+
```

A summary of the table without `WITH ROLLUP` looks like this:

```
mysql> SELECT name, size, SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size;
+------+-------+----------+
| name | size  | quantity |
+------+-------+----------+
| ball | small |       10 |
| ball | large |       20 |
| ball | NULL  |        5 |
| hoop | small |       15 |
| hoop | large |        5 |
| hoop | NULL  |        3 |
+------+-------+----------+
```

The result contains `NULL` values, but those do not represent super-aggregate rows because the query does not include `WITH ROLLUP`.

Adding `WITH ROLLUP` produces super-aggregate summary rows containing additional `NULL` values. However, without comparing this result to the previous one, it is not easy to see which `NULL` values occur in super-aggregate rows and which occur in regular grouped rows:

```
mysql> SELECT name, size, SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+------+-------+----------+
| name | size  | quantity |
+------+-------+----------+
| ball | NULL  |        5 |
| ball | large |       20 |
| ball | small |       10 |
| ball | NULL  |       35 |
| hoop | NULL  |        3 |
| hoop | large |        5 |
| hoop | small |       15 |
| hoop | NULL  |       23 |
| NULL | NULL  |       58 |
+------+-------+----------+
```

To distinguish `NULL` values in in super-aggregate rows from those in regular grouped rows, use `GROUPING()`, which returns 1 only for super-aggregate `NULL` values:

```
mysql> SELECT
         name, size, SUM(quantity) AS quantity,
         GROUPING(name) AS grp_name,
         GROUPING(size) AS grp_size
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+------+-------+----------+----------+----------+
| name | size  | quantity | grp_name | grp_size |
+------+-------+----------+----------+----------+
| ball | NULL  |        5 |        0 |        0 |
| ball | large |       20 |        0 |        0 |
| ball | small |       10 |        0 |        0 |
| ball | NULL  |       35 |        0 |        1 |
| hoop | NULL  |        3 |        0 |        0 |
```

```
| hoop | large |        5 |        0 |        0 |
| hoop | small |       15 |        0 |        0 |
| hoop | NULL  |       23 |        0 |        1 |
| NULL | NULL  |       58 |        1 |        1 |
+------+-------+----------+----------+----------+
```

Common uses for `GROUPING()`:

- Substitute a label for super-aggregate `NULL` values:

```
mysql> SELECT
         IF(GROUPING(name) = 1, 'All items', name) AS name,
         IF(GROUPING(size) = 1, 'All sizes', size) AS size,
         SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+-----------+-----------+----------+
| name      | size      | quantity |
+-----------+-----------+----------+
| ball      | NULL      |        5 |
| ball      | large     |       20 |
| ball      | small     |       10 |
| ball      | All sizes |       35 |
| hoop      | NULL      |        3 |
| hoop      | large     |        5 |
| hoop      | small     |       15 |
| hoop      | All sizes |       23 |
| All items | All sizes |       58 |
+-----------+-----------+----------+
```

- Return only super-aggregate lines by filtering out the regular grouped lines:

```
mysql> SELECT name, size, SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP
       HAVING GROUPING(name) = 1 OR GROUPING(size) = 1;
+------+------+----------+
| name | size | quantity |
+------+------+----------+
| ball | NULL |       35 |
| hoop | NULL |       23 |
| NULL | NULL |       58 |
+------+------+----------+
```

`GROUPING()` permits multiple expression arguments. In this case, the `GROUPING()` return value represents a bitmask combined from the results for each expression, where the lowest-order bit corresponds to the result for the rightmost expression. For example, with three expression arguments, `GROUPING(expr1, expr2, expr3)` is evaluated like this:

```
  result for GROUPING(expr3)
+ result for GROUPING(expr2) << 1
+ result for GROUPING(expr1) << 2
```

The following query shows how `GROUPING()` results for single arguments combine for a multiple-argument call to produce a bitmask value:

```
mysql> SELECT
         name, size, SUM(quantity) AS quantity,
         GROUPING(name) AS grp_name,
         GROUPING(size) AS grp_size,
       GROUPING(name, size) AS grp_all
       FROM t1
       GROUP BY name, size WITH ROLLUP;
+------+------+----------+----------+----------+---------+
| name | size | quantity | grp_name | grp_size | grp_all |
+------+------+----------+----------+----------+---------+
```

```
| ball | NULL  |        5 |        0 |        0 |       0 |
| ball | large |       20 |        0 |        0 |       0 |
| ball | small |       10 |        0 |        0 |       0 |
| ball | NULL  |       35 |        0 |        1 |       1 |
| hoop | NULL  |        3 |        0 |        0 |       0 |
| hoop | large |        5 |        0 |        0 |       0 |
| hoop | small |       15 |        0 |        0 |       0 |
| hoop | NULL  |       23 |        0 |        1 |       1 |
| NULL | NULL  |       58 |        1 |        1 |       3 |
+------+-------+----------+----------+----------+---------+
```

With multiple expression arguments, the `GROUPING()` return value is nonzero if any expression represents a super-aggregate value. Multiple-argument `GROUPING()` syntax thus provides a simpler way to write the earlier query that returned only super-aggregate rows, by using a single multiple-argument `GROUPING()` call rather than multiple single-argument calls:

```
mysql> SELECT name, size, SUM(quantity) AS quantity
       FROM t1
       GROUP BY name, size WITH ROLLUP
       HAVING GROUPING(name, size) <> 0;
+------+------+----------+
| name | size | quantity |
+------+------+----------+
| ball | NULL |       35 |
| hoop | NULL |       23 |
| NULL | NULL |       58 |
+------+------+----------+
```

Use of `GROUPING()` is subject to these limitations:

- Do not use subquery `GROUP BY` expressions as `GROUPING()` arguments because matching might fail. For example, matching fails for this query:

  ```
  mysql> SELECT GROUPING((SELECT MAX(name) FROM t1))
         FROM t1
         GROUP BY (SELECT MAX(name) FROM t1) WITH ROLLUP;
  ERROR 3580 (HY000): Argument #1 of GROUPING function is not in GROUP BY
  ```

- `GROUP BY` literal expressions should not be used within a `HAVING` clause as `GROUPING()` arguments. Due to differences between when the optimizer evaluates `GROUP BY` and `HAVING`, matching may succeed but `GROUPING()` evaluation does not produce the expected result. Consider this query:

  ```
  SELECT a AS f1, 'w' AS f2
  FROM t
  GROUP BY f1, f2 WITH ROLLUP
  HAVING GROUPING(f2) = 1;
  ```

  `GROUPING()` is evaluated earlier for the literal constant expression than for the `HAVING` clause as a whole and returns 0. To check whether a query such as this is affected, use `EXPLAIN` and look for `Impossible having` in the `Extra` column.

For more information about `WITH ROLLUP` and `GROUPING()`, see Section 12.18.2, “GROUP BY Modifiers”.

- `INET_ATON(expr)`

Given the dotted-quad representation of an IPv4 network address as a string, returns an integer that represents the numeric value of the address in network byte order (big endian). `INET_ATON()` returns `NULL` if it does not understand its argument.

```
mysql> SELECT INET_ATON('10.0.5.9');
       -> 167773449
```

For this example, the return value is calculated as $10\times256^3 + 0\times256^2 + 5\times256 + 9$.

`INET_ATON()` may or may not return a non-`NULL` result for short-form IP addresses (such as `'127.1'` as a representation of `'127.0.0.1'`). Because of this, `INET_ATON()`a should not be used for such addresses.

Note

To store values generated by `INET_ATON()`, use an `INT UNSIGNED` column rather than `INT`, which is signed. If you use a signed column, values corresponding to IP addresses for which the first octet is greater than 127 cannot be stored correctly. See Section 11.2.6, “Out-of-Range and Overflow Handling”.

- `INET_NTOA(`*`expr`*`)`

  Given a numeric IPv4 network address in network byte order, returns the dotted-quad string representation of the address as a nonbinary string in the connection character set. `INET_NTOA()` returns `NULL` if it does not understand its argument.

```
mysql> SELECT INET_NTOA(167773449);
        -> '10.0.5.9'
```

- `INET6_ATON(`*`expr`*`)`

  Given an IPv6 or IPv4 network address as a string, returns a binary string that represents the numeric value of the address in network byte order (big endian). Because numeric-format IPv6 addresses require more bytes than the largest integer type, the representation returned by this function has the `VARBINARY` data type: `VARBINARY(16)` for IPv6 addresses and `VARBINARY(4)` for IPv4 addresses. If the argument is not a valid address, `INET6_ATON()` returns `NULL`.

  The following examples use `HEX()` to display the `INET6_ATON()` result in printable form:

```
mysql> SELECT HEX(INET6_ATON('fdfe::5a55:caff:fefa:9089'));
        -> 'FDFE0000000000005A55CAFFFEFA9089'
mysql> SELECT HEX(INET6_ATON('10.0.5.9'));
        -> '0A000509'
```

  `INET6_ATON()` observes several constraints on valid arguments. These are given in the following list along with examples.

  - A trailing zone ID is not permitted, as in `fe80::3%1` or `fe80::3%eth0`.
  - A trailing network mask is not permitted, as in `2001:45f:3:ba::/64` or `192.168.1.0/24`.
  - For values representing IPv4 addresses, only classless addresses are supported. Classful addresses such as `192.168.1` are rejected. A trailing port number is not permitted, as in `192.168.1.2:8080`. Hexadecimal numbers in address components are not permitted, as in `192.0xa0.1.2`. Octal numbers are not supported: `192.168.010.1` is treated as `192.168.10.1`, not `192.168.8.1`. These IPv4 constraints also apply to IPv6 addresses that have IPv4 address parts, such as IPv4-compatible or IPv4-mapped addresses.

  To convert an IPv4 address *`expr`* represented in numeric form as an `INT` value to an IPv6 address represented in numeric form as a `VARBINARY` value, use this expression:

```
INET6_ATON(INET_NTOA(expr))
```

  For example:

```
mysql> SELECT HEX(INET6_ATON(INET_NTOA(167773449)));
       -> '0A000509'
```

- `INET6_NTOA(expr)`

  Given an IPv6 or IPv4 network address represented in numeric form as a binary string, returns the string representation of the address as a nonbinary string in the connection character set. If the argument is not a valid address, `INET6_NTOA()` returns `NULL`.

  `INET6_NTOA()` has these properties:

  - It does not use operating system functions to perform conversions, thus the output string is platform independent.
  - The return string has a maximum length of 39 (4 x 8 + 7). Given this statement:

    ```
    CREATE TABLE t AS SELECT INET6_NTOA(expr) AS c1;
    ```

    The resulting table would have this definition:

    ```
    CREATE TABLE t (c1 VARCHAR(39) CHARACTER SET utf8 DEFAULT NULL);
    ```

  - The return string uses lowercase letters for IPv6 addresses.

  ```
  mysql> SELECT INET6_NTOA(INET6_ATON('fdfe::5a55:caff:fefa:9089'));
         -> 'fdfe::5a55:caff:fefa:9089'
  mysql> SELECT INET6_NTOA(INET6_ATON('10.0.5.9'));
         -> '10.0.5.9'

  mysql> SELECT INET6_NTOA(UNHEX('FDFE0000000000005A55CAFFFEFA9089'));
         -> 'fdfe::5a55:caff:fefa:9089'
  mysql> SELECT INET6_NTOA(UNHEX('0A000509'));
         -> '10.0.5.9'
  ```

- `IS_FREE_LOCK(str)`

  Checks whether the lock named *`str`* is free to use (that is, not locked). Returns `1` if the lock is free (no one is using the lock), `0` if the lock is in use, and `NULL` if an error occurs (such as an incorrect argument).

  This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `IS_IPV4(expr)`

  Returns 1 if the argument is a valid IPv4 address specified as a string, 0 otherwise.

  ```
  mysql> SELECT IS_IPV4('10.0.5.9'), IS_IPV4('10.0.5.256');
         -> 1, 0
  ```

  For a given argument, if `IS_IPV4()` returns 1, `INET_ATON()` (and `INET6_ATON()`) will return non-`NULL`. The converse statement is not true: In some cases, `INET_ATON()` returns non-`NULL` when `IS_IPV4()` returns 0.

  As implied by the preceding remarks, `IS_IPV4()` is more strict than `INET_ATON()` about what constitutes a valid IPv4 address, so it may be useful for applications that need to perform strong checks against invalid values. Alternatively, use `INET6_ATON()` to convert IPv4 addresses to internal form and check for a `NULL` result (which indicates an invalid address). `INET6_ATON()` is equally strong as `IS_IPV4()` about checking IPv4 addresses.

- `IS_IPV4_COMPAT(expr)`

  This function takes an IPv6 address represented in numeric form as a binary string, as returned by `INET6_ATON()`. It returns 1 if the argument is a valid IPv4-compatible IPv6 address, 0 otherwise. IPv4-compatible addresses have the form `::ipv4_address`.

```
mysql> SELECT IS_IPV4_COMPAT(INET6_ATON('::10.0.5.9'));
        -> 1
mysql> SELECT IS_IPV4_COMPAT(INET6_ATON('::ffff:10.0.5.9'));
        -> 0
```

  The IPv4 part of an IPv4-compatible address can also be represented using hexadecimal notation. For example, `192.168.0.1` has this raw hexadecimal value:

```
mysql> SELECT HEX(INET6_ATON('192.168.0.1'));
        -> 'C0A80001'
```

  Expressed in IPv4-compatible form, `::192.168.0.1` is equivalent to `::c0a8:0001` or (without leading zeros) `::c0a8:1`

```
mysql> SELECT
    ->   IS_IPV4_COMPAT(INET6_ATON('::192.168.0.1')),
    ->   IS_IPV4_COMPAT(INET6_ATON('::c0a8:0001')),
    ->   IS_IPV4_COMPAT(INET6_ATON('::c0a8:1'));
        -> 1, 1, 1
```

- `IS_IPV4_MAPPED(expr)`

  This function takes an IPv6 address represented in numeric form as a binary string, as returned by `INET6_ATON()`. It returns 1 if the argument is a valid IPv4-mapped IPv6 address, 0 otherwise. IPv4-mapped addresses have the form `::ffff:ipv4_address`.

```
mysql> SELECT IS_IPV4_MAPPED(INET6_ATON('::10.0.5.9'));
        -> 0
mysql> SELECT IS_IPV4_MAPPED(INET6_ATON('::ffff:10.0.5.9'));
        -> 1
```

  As with `IS_IPV4_COMPAT()` the IPv4 part of an IPv4-mapped address can also be represented using hexadecimal notation:

```
mysql> SELECT
    ->   IS_IPV4_MAPPED(INET6_ATON('::ffff:192.168.0.1')),
    ->   IS_IPV4_MAPPED(INET6_ATON('::ffff:c0a8:0001')),
    ->   IS_IPV4_MAPPED(INET6_ATON('::ffff:c0a8:1'));
        -> 1, 1, 1
```

- `IS_IPV6(expr)`

  Returns 1 if the argument is a valid IPv6 address specified as a string, 0 otherwise. This function does not consider IPv4 addresses to be valid IPv6 addresses.

```
mysql> SELECT IS_IPV6('10.0.5.9'), IS_IPV6('::1');
        -> 0, 1
```

  For a given argument, if `IS_IPV6()` returns 1, `INET6_ATON()` will return non-`NULL`.

- `IS_USED_LOCK(str)`

  Checks whether the lock named `str` is in use (that is, locked). If so, it returns the connection identifier of the client session that holds the lock. Otherwise, it returns `NULL`.

This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `IS_UUID(string_uuid)`

  Returns 1 if the argument is a valid string-format UUID, 0 if the argument is not a valid UUID, and `NULL` if the argument is `NULL`.

  "Valid" means that the value is in a format that can be parsed. That is, it has the correct length and contains only the permitted characters (hexadecimal digits in any lettercase and, optionally, dashes and curly braces). This format is most common:

```
aaaaaaaa-bbbb-cccc-dddd-eeeeeeeeeeee
```

  These other formats are also permitted:

```
aaaaaaaabbbbccccddddeeeeeeeeeeee
{aaaaaaaa-bbbb-cccc-dddd-eeeeeeeeeeee}
```

  For the meanings of fields within the value, see the `UUID()` function description.

```
mysql> SELECT IS_UUID('6ccd780c-baba-1026-9564-5b8c656024db');
+------------------------------------------------+
| IS_UUID('6ccd780c-baba-1026-9564-5b8c656024db') |
+------------------------------------------------+
|                                              1 |
+------------------------------------------------+
mysql> SELECT IS_UUID('6CCD780C-BABA-1026-9564-5B8C656024DB');
+------------------------------------------------+
| IS_UUID('6CCD780C-BABA-1026-9564-5B8C656024DB') |
+------------------------------------------------+
|                                              1 |
+------------------------------------------------+
mysql> SELECT IS_UUID('6ccd780cbaba102695645b8c656024db');
+--------------------------------------------+
| IS_UUID('6ccd780cbaba102695645b8c656024db') |
+--------------------------------------------+
|                                          1 |
+--------------------------------------------+
mysql> SELECT IS_UUID('{6ccd780c-baba-1026-9564-5b8c656024db}');
+--------------------------------------------------+
| IS_UUID('{6ccd780c-baba-1026-9564-5b8c656024db}') |
+--------------------------------------------------+
|                                                1 |
+--------------------------------------------------+
mysql> SELECT IS_UUID('6ccd780c-baba-1026-9564-5b8c6560');
+--------------------------------------------+
| IS_UUID('6ccd780c-baba-1026-9564-5b8c6560') |
+--------------------------------------------+
|                                          0 |
+--------------------------------------------+
mysql> SELECT IS_UUID(RAND());
+-----------------+
| IS_UUID(RAND()) |
+-----------------+
|               0 |
+-----------------+
```

- `MASTER_POS_WAIT(log_name,log_pos[,timeout][,channel_name])`

  This function is useful for control of master/slave synchronization. It blocks until the slave has read and applied all updates up to the specified position in the master log. The return value is the number of log events the slave had to wait for to advance to the specified position. The function returns `NULL` if the slave SQL thread is not started, the slave's master information is not initialized, the arguments

are incorrect, or an error occurs. It returns `-1` if the timeout has been exceeded. If the slave SQL thread stops while `MASTER_POS_WAIT()` is waiting, the function returns `NULL`. If the slave is past the specified position, the function returns immediately.

If a *`timeout`* value is specified, `MASTER_POS_WAIT()` stops waiting when *`timeout`* seconds have elapsed. *`timeout`* must be greater than 0; a zero or negative *`timeout`* means no timeout.

The optional *`channel`* a enables you to choose which replication channel the function applies to. See Section 18.2.3, “Replication Channels” for more information.

This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `NAME_CONST(`*`name`*`,`*`value`*`)`

  Returns the given value. When used to produce a result set column, `NAME_CONST()` causes the column to have the given name. The arguments should be constants.

```
mysql> SELECT NAME_CONST('myname', 14);
+--------+
| myname |
+--------+
|     14 |
+--------+
```

  This function is for internal use only. The server uses it when writing statements from stored programs that contain references to local program variables, as described in Section 23.7, “Binary Logging of Stored Programs”, You might see this function in the output from `mysqlbinlog`.

  For your applications, you can obtain exactly the same result as in the example just shown by using simple aliasing, like this:

```
mysql> SELECT 14 AS myname;
+--------+
| myname |
+--------+
|     14 |
+--------+
1 row in set (0.00 sec)
```

  See Section 13.2.10, “SELECT Syntax”, for more information about column aliases.

- `RELEASE_ALL_LOCKS()`

  Releases all named locks held by the current session and returns the number of locks released (0 if there were none)

  This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `RELEASE_LOCK(`*`str`*`)`

  Releases the lock named by the string *`str`* that was obtained with `GET_LOCK()`. Returns `1` if the lock was released, `0` if the lock was not established by this thread (in which case the lock is not released), and `NULL` if the named lock did not exist. The lock does not exist if it was never obtained by a call to `GET_LOCK()` or if it has previously been released.

  The `DO` statement is convenient to use with `RELEASE_LOCK()`. See Section 13.2.3, “DO Syntax”.

  This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `SLEEP(duration)`

  Sleeps (pauses) for the number of seconds given by the *`duration`* argument, then returns 0. The duration may have a fractional part. If the argument is `NULL` or negative, `SLEEP()` produces a warning, or an error in strict SQL mode.

  When sleep returns normally (without interruption), it returns 0:

```
mysql> SELECT SLEEP(1000);
+-------------+
| SLEEP(1000) |
+-------------+
|           0 |
+-------------+
```

  When `SLEEP()` is the only thing invoked by a query that is interrupted, it returns 1 and the query itself returns no error. This is true whether the query is killed or times out:

  - This statement is interrupted using `KILL QUERY` from another session:

```
mysql> SELECT SLEEP(1000);
+-------------+
| SLEEP(1000) |
+-------------+
|           1 |
+-------------+
```

  - This statement is interrupted by timing out:

```
mysql> SELECT /*+ MAX_EXECUTION_TIME(1) */ SLEEP(1000);
+-------------+
| SLEEP(1000) |
+-------------+
|           1 |
+-------------+
```

  When `SLEEP()` is only part of a query that is interrupted, the query returns an error:

  - This statement is interrupted using `KILL QUERY` from another session:

```
mysql> SELECT 1 FROM t1 WHERE SLEEP(1000);
ERROR 1317 (70100): Query execution was interrupted
```

  - This statement is interrupted by timing out:

```
mysql> SELECT /*+ MAX_EXECUTION_TIME(1000) */ 1 FROM t1 WHERE SLEEP(1000);
ERROR 3024 (HY000): Query execution was interrupted, maximum statement
execution time exceeded
```

  This function is unsafe for statement-based replication. A warning is logged if you use this function when `binlog_format` is set to `STATEMENT`.

- `UUID()`

  Returns a Universal Unique Identifier (UUID) generated according to RFC 4122, “A Universally Unique IDentifier (UUID) URN Namespace” (http://www.ietf.org/rfc/rfc4122.txt).

  A UUID is designed as a number that is globally unique in space and time. Two calls to `UUID()` are expected to generate two different values, even if these calls are performed on two separate devices not connected to each other.

Warning

Although `UUID()` values are intended to be unique, they are not necessarily unguessable or unpredictable. If unpredictability is required, UUID values should be generated some other way.

`UUID()` returns a value that conforms to UUID version 1 as described in RFC 4122. The value is a 128-bit number represented as a `utf8` string of five hexadecimal numbers in `aaaaaaaa-bbbb-cccc-dddd-eeeeeeeeeeee` format:

- The first three numbers are generated from the low, middle, and high parts of a timestamp. The high part also includes the UUID version number.
- The fourth number preserves temporal uniqueness in case the timestamp value loses monotonicity (for example, due to daylight saving time).
- The fifth number is an IEEE 802 node number that provides spatial uniqueness. A random number is substituted if the latter is not available (for example, because the host device has no Ethernet card, or it is unknown how to find the hardware address of an interface on the host operating system). In this case, spatial uniqueness cannot be guaranteed. Nevertheless, a collision should have *very* low probability.

  The MAC address of an interface is taken into account only on FreeBSD and Linux. On other operating systems, MySQL uses a randomly generated 48-bit number.

```
mysql> SELECT UUID();
        -> '6ccd780c-baba-1026-9564-5b8c656024db'
```

To convert between string and binary UUID values, use the `UUID_TO_BIN()` and `BIN_TO_UUID()` functions. To check whether a string is a valid UUID value, use the `IS_UUID()` function.

Note

`UUID()` does not work with statement-based replication.

- `UUID_SHORT()`

  Returns a “short” universal identifier as a 64-bit unsigned integer. Values returned by `UUID_SHORT()` differ from the string-format 128-bit identifiers returned by the `UUID()` function and have different uniqueness properties. The value of `UUID_SHORT()` is guaranteed to be unique if the following conditions hold:

  - The `server_id` value of the current server is between 0 and 255 and is unique among your set of master and slave servers
  - You do not set back the system time for your server host between `mysqld` restarts
  - You invoke `UUID_SHORT()` on average fewer than 16 million times per second between `mysqld` restarts

  The `UUID_SHORT()` return value is constructed this way:

```
  (server_id & 255) << 56
+ (server_startup_time_in_seconds << 24)
+ incremented_variable++;
```

```
mysql> SELECT UUID_SHORT();
        -> 92395783831158784
```

Note

`UUID_SHORT()` does not work with statement-based replication.

- `UUID_TO_BIN(string_uuid)`, `UUID_TO_BIN(string_uuid, swap_flag)`

  Converts a string UUID to a binary UUID and returns the result. (The `IS_UUID()` function description lists the permitted string UUID formats.) The return binary UUID is a `VARBINARY(16)` value. If the UUID argument is `NULL`, the return value is `NULL`. If any argument is invalid, an error occurs.

  `UUID_TO_BIN()` takes one or two arguments:

  - The one-argument form takes a string UUID value. The binary result is in the same order as the string argument.
  - The two-argument form takes a string UUID value and a flag value:
    - If *`swap_flag`* is 0, the two-argument form is equivalent to the one-argument form. The binary result is in the same order as the string argument.
    - If *`swap_flag`* is 1, the format of the return value differs: The time-low and time-high parts (the first and third groups of hexadecimal digits, respectively) are swapped. This moves the more rapidly varying part to the right and can improve indexing efficiency if the result is stored in an indexed column.

  Time-part swapping assumes the use of UUID version 1 values, such as are generated by the `UUID()` function. For UUID values produced by other means that do not follow version 1 format, time-part swapping provides no benefit. For details about version 1 format, see the `UUID()` function description.

  Suppose that you have the following string UUID value:

```
mysql> SET @uuid = '6ccd780c-baba-1026-9564-5b8c656024db';
```

  To convert the string UUID to binary with or without time-part swapping, use `UUID_TO_BIN()`:

```
mysql> SELECT HEX(UUID_TO_BIN(@uuid));
+----------------------------------+
| HEX(UUID_TO_BIN(@uuid))          |
+----------------------------------+
| 6CCD780CBABA102695645B8C656024DB |
+----------------------------------+
mysql> SELECT HEX(UUID_TO_BIN(@uuid, 0));
+----------------------------------+
| HEX(UUID_TO_BIN(@uuid, 0))       |
+----------------------------------+
| 6CCD780CBABA102695645B8C656024DB |
+----------------------------------+
mysql> SELECT HEX(UUID_TO_BIN(@uuid, 1));
+----------------------------------+
| HEX(UUID_TO_BIN(@uuid, 1))       |
+----------------------------------+
| 1026BABA6CCD780C95645B8C656024DB |
+----------------------------------+
```

  To convert a binary UUID returned by `UUID_TO_BIN()` to a string UUID, use `BIN_TO_UUID()`. If you produce a binary UUID by calling `UUID_TO_BIN()` with a second argument of 1 to swap time parts, you should also pass a second argument of 1 to `BIN_TO_UUID()` to unswap the time parts when converting the binary UUID back to a string UUID:

```
mysql> SELECT BIN_TO_UUID(UUID_TO_BIN(@uuid));
+--------------------------------------+
| BIN_TO_UUID(UUID_TO_BIN(@uuid))      |
+--------------------------------------+
| 6ccd780c-baba-1026-9564-5b8c656024db |
+--------------------------------------+
mysql> SELECT BIN_TO_UUID(UUID_TO_BIN(@uuid,0),0);
+--------------------------------------+
| BIN_TO_UUID(UUID_TO_BIN(@uuid,0),0)  |
+--------------------------------------+
| 6ccd780c-baba-1026-9564-5b8c656024db |
+--------------------------------------+
mysql> SELECT BIN_TO_UUID(UUID_TO_BIN(@uuid,1),1);
+--------------------------------------+
| BIN_TO_UUID(UUID_TO_BIN(@uuid,1),1)  |
+--------------------------------------+
| 6ccd780c-baba-1026-9564-5b8c656024db |
+--------------------------------------+
```

If the use of time-part swapping is not the same for the conversion in both directions, the original UUID will not be recovered properly:

```
mysql> SELECT BIN_TO_UUID(UUID_TO_BIN(@uuid,0),1);
+--------------------------------------+
| BIN_TO_UUID(UUID_TO_BIN(@uuid,0),1)  |
+--------------------------------------+
| baba1026-780c-6ccd-9564-5b8c656024db |
+--------------------------------------+
mysql> SELECT BIN_TO_UUID(UUID_TO_BIN(@uuid,1),0);
+--------------------------------------+
| BIN_TO_UUID(UUID_TO_BIN(@uuid,1),0)  |
+--------------------------------------+
| 1026baba-6ccd-780c-9564-5b8c656024db |
+--------------------------------------+
```

- `VALUES(col_name)`

  In an `INSERT ... ON DUPLICATE KEY UPDATE` statement, you can use the `VALUES(col_name)` function in the `UPDATE` clause to refer to column values from the `INSERT` portion of the statement. In other words, `VALUES(col_name)` in the `UPDATE` clause refers to the value of `col_name` that would be inserted, had no duplicate-key conflict occurred. This function is especially useful in multiple-row inserts. The `VALUES()` function is meaningful only in the `ON DUPLICATE KEY UPDATE` clause of `INSERT` statements and returns `NULL` otherwise. See Section 13.2.6.3, “INSERT ... ON DUPLICATE KEY UPDATE Syntax”.

```
mysql> INSERT INTO table (a,b,c) VALUES (1,2,3),(4,5,6)
    -> ON DUPLICATE KEY UPDATE c=VALUES(a)+VALUES(b);
```

12.21 Precision Math

MySQL provides support for precision math: numeric value handling that results in extremely accurate results and a high degree control over invalid values. Precision math is based on these two features:

- SQL modes that control how strict the server is about accepting or rejecting invalid data.
- The MySQL library for fixed-point arithmetic.

These features have several implications for numeric operations and provide a high degree of compliance with standard SQL:

- **Precise calculations**: For exact-value numbers, calculations do not introduce floating-point errors. Instead, exact precision is used. For example, MySQL treats a number such as `.0001` as an exact

value rather than as an approximation, and summing it 10,000 times produces a result of exactly `1`, not a value that is merely “close” to 1.

- **Well-defined rounding behavior**: For exact-value numbers, the result of `ROUND()` depends on its argument, not on environmental factors such as how the underlying C library works.
- **Platform independence**: Operations on exact numeric values are the same across different platforms such as Windows and Unix.
- **Control over handling of invalid values**: Overflow and division by zero are detectable and can be treated as errors. For example, you can treat a value that is too large for a column as an error rather than having the value truncated to lie within the range of the column's data type. Similarly, you can treat division by zero as an error rather than as an operation that produces a result of `NULL`. The choice of which approach to take is determined by the setting of the server SQL mode.

The following discussion covers several aspects of how precision math works, including possible incompatibilities with older applications. At the end, some examples are given that demonstrate how MySQL handles numeric operations precisely. For information about controlling the SQL mode, see Section 5.1.8, “Server SQL Modes”.

12.21.1 Types of Numeric Values

The scope of precision math for exact-value operations includes the exact-value data types (integer and `DECIMAL` types) and exact-value numeric literals. Approximate-value data types and numeric literals are handled as floating-point numbers.

Exact-value numeric literals have an integer part or fractional part, or both. They may be signed. Examples: `1`, `.2`, `3.4`, `-5`, `-6.78`, `+9.10`.

Approximate-value numeric literals are represented in scientific notation with a mantissa and exponent. Either or both parts may be signed. Examples: `1.2E3`, `1.2E-3`, `-1.2E3`, `-1.2E-3`.

Two numbers that look similar may be treated differently. For example, `2.34` is an exact-value (fixed-point) number, whereas `2.34E0` is an approximate-value (floating-point) number.

The `DECIMAL` data type is a fixed-point type and calculations are exact. In MySQL, the `DECIMAL` type has several synonyms: `NUMERIC`, `DEC`, `FIXED`. The integer types also are exact-value types.

The `FLOAT` and `DOUBLE` data types are floating-point types and calculations are approximate. In MySQL, types that are synonymous with `FLOAT` or `DOUBLE` are `DOUBLE PRECISION` and `REAL`.

12.21.2 DECIMAL Data Type Characteristics

This section discusses the characteristics of the `DECIMAL` data type (and its synonyms), with particular regard to the following topics:

- Maximum number of digits
- Storage format
- Storage requirements
- The nonstandard MySQL extension to the upper range of `DECIMAL` columns

The declaration syntax for a `DECIMAL` column is `DECIMAL(M,D)`. The ranges of values for the arguments are as follows:

- *`M`* is the maximum number of digits (the precision). It has a range of 1 to 65.
- *`D`* is the number of digits to the right of the decimal point (the scale). It has a range of 0 to 30 and must be no larger than *`M`*.

The maximum value of 65 for *M* means that calculations on DECIMAL values are accurate up to 65 digits. This limit of 65 digits of precision also applies to exact-value numeric literals, so the maximum range of such literals differs from before.

Values for DECIMAL columns are stored using a binary format that packs nine decimal digits into 4 bytes. The storage requirements for the integer and fractional parts of each value are determined separately. Each multiple of nine digits requires 4 bytes, and any remaining digits left over require some fraction of 4 bytes. The storage required for remaining digits is given by the following table.

Leftover Digits	Number of Bytes
0	0
1–2	1
3–4	2
5–6	3
7–9	4

For example, a DECIMAL(18,9) column has nine digits on either side of the decimal point, so the integer part and the fractional part each require 4 bytes. A DECIMAL(20,6) column has fourteen integer digits and six fractional digits. The integer digits require four bytes for nine of the digits and 3 bytes for the remaining five digits. The six fractional digits require 3 bytes.

DECIMAL columns do not store a leading + character or - character or leading 0 digits. If you insert +0003.1 into a DECIMAL(5,1) column, it is stored as 3.1. For negative numbers, a literal - character is not stored.

DECIMAL columns do not permit values larger than the range implied by the column definition. For example, a DECIMAL(3,0) column supports a range of -999 to 999. A DECIMAL(*M*,*D*) column permits at most *M* - *D* digits to the left of the decimal point.

The SQL standard requires that the precision of NUMERIC(*M*,*D*) be *exactly* *M* digits. For DECIMAL(*M*,*D*), the standard requires a precision of at least *M* digits but permits more. In MySQL, DECIMAL(*M*,*D*) and NUMERIC(*M*,*D*) are the same, and both have a precision of exactly *M* digits.

For a full explanation of the internal format of DECIMAL values, see the file strings/decimal.c in a MySQL source distribution. The format is explained (with an example) in the decimal2bin() function.

12.21.3 Expression Handling

With precision math, exact-value numbers are used as given whenever possible. For example, numbers in comparisons are used exactly as given without a change in value. In strict SQL mode, for INSERT into a column with an exact data type (DECIMAL or integer), a number is inserted with its exact value if it is within the column range. When retrieved, the value should be the same as what was inserted. (If strict SQL mode is not enabled, truncation for INSERT is permissible.)

Handling of a numeric expression depends on what kind of values the expression contains:

- If any approximate values are present, the expression is approximate and is evaluated using floating-point arithmetic.
- If no approximate values are present, the expression contains only exact values. If any exact value contains a fractional part (a value following the decimal point), the expression is evaluated using DECIMAL exact arithmetic and has a precision of 65 digits. The term “exact” is subject to the limits of what can be represented in binary. For example, 1.0/3.0 can be approximated in decimal notation as .333..., but not written as an exact number, so (1.0/3.0)*3.0 does not evaluate to exactly 1.0.
- Otherwise, the expression contains only integer values. The expression is exact and is evaluated using integer arithmetic and has a precision the same as BIGINT (64 bits).

If a numeric expression contains any strings, they are converted to double-precision floating-point values and the expression is approximate.

Inserts into numeric columns are affected by the SQL mode, which is controlled by the `sql_mode` system variable. (See Section 5.1.8, “Server SQL Modes”.) The following discussion mentions strict mode (selected by the `STRICT_ALL_TABLES` or `STRICT_TRANS_TABLES` mode values) and `ERROR_FOR_DIVISION_BY_ZERO`. To turn on all restrictions, you can simply use `TRADITIONAL` mode, which includes both strict mode values and `ERROR_FOR_DIVISION_BY_ZERO`:

```
mysql> SET sql_mode='TRADITIONAL';
```

If a number is inserted into an exact type column (`DECIMAL` or integer), it is inserted with its exact value if it is within the column range.

If the value has too many digits in the fractional part, rounding occurs and a warning is generated. Rounding is done as described in Section 12.21.4, “Rounding Behavior”.

If the value has too many digits in the integer part, it is too large and is handled as follows:

- If strict mode is not enabled, the value is truncated to the nearest legal value and a warning is generated.
- If strict mode is enabled, an overflow error occurs.

Underflow is not detected, so underflow handling is undefined.

For inserts of strings into numeric columns, conversion from string to number is handled as follows if the string has nonnumeric contents:

- A string that does not begin with a number cannot be used as a number and produces an error in strict mode, or a warning otherwise. This includes the empty string.
- A string that begins with a number can be converted, but the trailing nonnumeric portion is truncated. If the truncated portion contains anything other than spaces, this produces an error in strict mode, or a warning otherwise.

By default, division by zero produces a result of `NULL` and no warning. By setting the SQL mode appropriately, division by zero can be restricted.

With the `ERROR_FOR_DIVISION_BY_ZERO` SQL mode enabled, MySQL handles division by zero differently:

- If strict mode is not enabled, a warning occurs.
- If strict mode is enabled, inserts and updates involving division by zero are prohibited, and an error occurs.

In other words, inserts and updates involving expressions that perform division by zero can be treated as errors, but this requires `ERROR_FOR_DIVISION_BY_ZERO` in addition to strict mode.

Suppose that we have this statement:

```
INSERT INTO t SET i = 1/0;
```

This is what happens for combinations of strict and `ERROR_FOR_DIVISION_BY_ZERO` modes.

`sql_mode` Value	Result
`''` (Default)	No warning, no error; `i` is set to `NULL`.
strict	No warning, no error; `i` is set to `NULL`.
`ERROR_FOR_DIVISION_BY_ZERO`	Warning, no error; `i` is set to `NULL`.

sql_mode Value	Result
strict,ERROR_FOR_DIVISION_BY_ZERO	Error condition; no row is inserted.

12.21.4 Rounding Behavior

This section discusses precision math rounding for the ROUND() function and for inserts into columns with exact-value types (DECIMAL and integer).

The ROUND() function rounds differently depending on whether its argument is exact or approximate:

- For exact-value numbers, ROUND() uses the “round half up” rule: A value with a fractional part of .5 or greater is rounded up to the next integer if positive or down to the next integer if negative. (In other words, it is rounded away from zero.) A value with a fractional part less than .5 is rounded down to the next integer if positive or up to the next integer if negative.
- For approximate-value numbers, the result depends on the C library. On many systems, this means that ROUND() uses the “round to nearest even” rule: A value with any fractional part is rounded to the nearest even integer.

The following example shows how rounding differs for exact and approximate values:

```
mysql> SELECT ROUND(2.5), ROUND(25E-1);
+------------+--------------+
| ROUND(2.5) | ROUND(25E-1) |
+------------+--------------+
| 3          |            2 |
+------------+--------------+
```

For inserts into a DECIMAL or integer column, the target is an exact data type, so rounding uses “round half away from zero,” regardless of whether the value to be inserted is exact or approximate:

```
mysql> CREATE TABLE t (d DECIMAL(10,0));
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t VALUES(2.5),(2.5E0);
Query OK, 2 rows affected, 2 warnings (0.00 sec)
Records: 2  Duplicates: 0  Warnings: 2

mysql> SELECT d FROM t;
+------+
| d    |
+------+
| 3    |
| 3    |
+------+
```

12.21.5 Precision Math Examples

This section provides some examples that show precision math query results in MySQL. These examples demonstrate the principles described in Section 12.21.3, “Expression Handling”, and Section 12.21.4, “Rounding Behavior”.

Example 1. Numbers are used with their exact value as given when possible:

```
mysql> SELECT (.1 + .2) = .3;
+----------------+
| (.1 + .2) = .3 |
+----------------+
|              1 |
+----------------+
```

For floating-point values, results are inexact:

```
mysql> SELECT (.1E0 + .2E0) = .3E0;
+----------------------+
| (.1E0 + .2E0) = .3E0 |
+----------------------+
|                    0 |
+----------------------+
```

Another way to see the difference in exact and approximate value handling is to add a small number to a sum many times. Consider the following stored procedure, which adds `.0001` to a variable 1,000 times.

```
CREATE PROCEDURE p ()
BEGIN
  DECLARE i INT DEFAULT 0;
  DECLARE d DECIMAL(10,4) DEFAULT 0;
  DECLARE f FLOAT DEFAULT 0;
  WHILE i < 10000 DO
    SET d = d + .0001;
    SET f = f + .0001E0;
    SET i = i + 1;
  END WHILE;
  SELECT d, f;
END;
```

The sum for both `d` and `f` logically should be 1, but that is true only for the decimal calculation. The floating-point calculation introduces small errors:

```
+--------+------------------+
| d      | f                |
+--------+------------------+
| 1.0000 | 0.99999999999991 |
+--------+------------------+
```

Example 2. Multiplication is performed with the scale required by standard SQL. That is, for two numbers *`X1`* and *`X2`* that have scale *`S1`* and *`S2`*, the scale of the result is *`S1 + S2`*:

```
mysql> SELECT .01 * .01;
+-----------+
| .01 * .01 |
+-----------+
| 0.0001    |
+-----------+
```

Example 3. Rounding behavior for exact-value numbers is well-defined:

Rounding behavior (for example, with the `ROUND()` function) is independent of the implementation of the underlying C library, which means that results are consistent from platform to platform.

- Rounding for exact-value columns (`DECIMAL` and integer) and exact-valued numbers uses the “round half away from zero” rule. Values with a fractional part of .5 or greater are rounded away from zero to the nearest integer, as shown here:

  ```
  mysql> SELECT ROUND(2.5), ROUND(-2.5);
  +------------+-------------+
  | ROUND(2.5) | ROUND(-2.5) |
  +------------+-------------+
  | 3          | -3          |
  +------------+-------------+
  ```

- Rounding for floating-point values uses the C library, which on many systems uses the “round to nearest even” rule. Values with any fractional part on such systems are rounded to the nearest even integer:

```
mysql> SELECT ROUND(2.5E0), ROUND(-2.5E0);
+--------------+---------------+
| ROUND(2.5E0) | ROUND(-2.5E0) |
+--------------+---------------+
|            2 |            -2 |
+--------------+---------------+
```

Example 4. In strict mode, inserting a value that is out of range for a column causes an error, rather than truncation to a legal value.

When MySQL is not running in strict mode, truncation to a legal value occurs:

```
mysql> SET sql_mode='';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t (i TINYINT);
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO t SET i = 128;
Query OK, 1 row affected, 1 warning (0.00 sec)

mysql> SELECT i FROM t;
+------+
| i    |
+------+
|  127 |
+------+
1 row in set (0.00 sec)
```

However, an error occurs if strict mode is in effect:

```
mysql> SET sql_mode='STRICT_ALL_TABLES';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t (i TINYINT);
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t SET i = 128;
ERROR 1264 (22003): Out of range value adjusted for column 'i' at row 1

mysql> SELECT i FROM t;
Empty set (0.00 sec)
```

Example 5: In strict mode and with `ERROR_FOR_DIVISION_BY_ZERO` set, division by zero causes an error, not a result of `NULL`.

In nonstrict mode, division by zero has a result of `NULL`:

```
mysql> SET sql_mode='';
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t (i TINYINT);
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t SET i = 1 / 0;
Query OK, 1 row affected (0.00 sec)

mysql> SELECT i FROM t;
+------+
| i    |
+------+
| NULL |
+------+
1 row in set (0.03 sec)
```

However, division by zero is an error if the proper SQL modes are in effect:

```
mysql> SET sql_mode='STRICT_ALL_TABLES,ERROR_FOR_DIVISION_BY_ZERO';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t (i TINYINT);
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO t SET i = 1 / 0;
ERROR 1365 (22012): Division by 0

mysql> SELECT i FROM t;
Empty set (0.01 sec)
```

Example 6. Exact-value literals are evaluated as exact values.

Approximate-value literals are evaluated using floating point, but exact-value literals are handled as `DECIMAL`:

```
mysql> CREATE TABLE t SELECT 2.5 AS a, 25E-1 AS b;
Query OK, 1 row affected (0.01 sec)
Records: 1  Duplicates: 0  Warnings: 0

mysql> DESCRIBE t;
+-------+----------------------+------+-----+---------+-------+
| Field | Type                 | Null | Key | Default | Extra |
+-------+----------------------+------+-----+---------+-------+
| a     | decimal(2,1) unsigned | NO   |     | 0.0     |       |
| b     | double               | NO   |     | 0       |       |
+-------+----------------------+------+-----+---------+-------+
2 rows in set (0.01 sec)
```

Example 7. If the argument to an aggregate function is an exact numeric type, the result is also an exact numeric type, with a scale at least that of the argument.

Consider these statements:

```
mysql> CREATE TABLE t (i INT, d DECIMAL, f FLOAT);
mysql> INSERT INTO t VALUES(1,1,1);
mysql> CREATE TABLE y SELECT AVG(i), AVG(d), AVG(f) FROM t;
```

The result is a double only for the floating-point argument. For exact type arguments, the result is also an exact type:

```
mysql> DESCRIBE y;
+--------+--------------+------+-----+---------+-------+
| Field  | Type         | Null | Key | Default | Extra |
+--------+--------------+------+-----+---------+-------+
| AVG(i) | decimal(14,4) | YES  |     | NULL    |       |
| AVG(d) | decimal(14,4) | YES  |     | NULL    |       |
| AVG(f) | double       | YES  |     | NULL    |       |
+--------+--------------+------+-----+---------+-------+
```

The result is a double only for the floating-point argument. For exact type arguments, the result is also an exact type.

Chapter 13 SQL Statement Syntax

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This chapter describes the syntax for the SQL statements supported by MySQL.

13.1 Data Definition Statements

13.1.1 ALTER DATABASE Syntax

```
ALTER {DATABASE | SCHEMA} [db_name]
    alter_specification ...

alter_specification:
    [DEFAULT] CHARACTER SET [=] charset_name
  | [DEFAULT] COLLATE [=] collation_name
```

`ALTER DATABASE` enables you to change the overall characteristics of a database. These characteristics are stored in the data dictionary. To use `ALTER DATABASE`, you need the `ALTER` privilege on the database. `ALTER SCHEMA` is a synonym for `ALTER DATABASE`.

The database name can be omitted from the first syntax, in which case the statement applies to the default database.

National Language Characteristics

The `CHARACTER SET` clause changes the default database character set. The `COLLATE` clause changes the default database collation. Section 10.1, “Character Set Support”, discusses character set and collation names.

You can see what character sets and collations are available using, respectively, the `SHOW CHARACTER SET` and `SHOW COLLATION` statements. See Section 13.7.5.3, “SHOW CHARACTER SET Syntax”, and Section 13.7.5.4, “SHOW COLLATION Syntax”, for more information.

If you change the default character set or collation for a database, stored routines that use the database defaults must be dropped and recreated so that they use the new defaults. (In a stored routine, variables with character data types use the database defaults if the character set or collation are not specified explicitly. See Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”.)

13.1.2 ALTER EVENT Syntax

```
ALTER
    [DEFINER = { user | CURRENT_USER }]
    EVENT event_name
    [ON SCHEDULE schedule]
    [ON COMPLETION [NOT] PRESERVE]
    [RENAME TO new_event_name]
    [ENABLE | DISABLE | DISABLE ON SLAVE]
    [COMMENT 'comment']
    [DO event_body]
```

The `ALTER EVENT` statement changes one or more of the characteristics of an existing event without the need to drop and recreate it. The syntax for each of the `DEFINER`, `ON SCHEDULE`, `ON COMPLETION`, `COMMENT`, `ENABLE` / `DISABLE`, and `DO` clauses is exactly the same as when used with `CREATE EVENT`. (See Section 13.1.10, “CREATE EVENT Syntax”.)

Any user can alter an event defined on a database for which that user has the `EVENT` privilege. When a user executes a successful `ALTER EVENT` statement, that user becomes the definer for the affected event.

`ALTER EVENT` works only with an existing event:

```
mysql> ALTER EVENT no_such_event
     >     ON SCHEDULE
     >       EVERY '2:3' DAY_HOUR;
ERROR 1517 (HY000): Unknown event 'no_such_event'
```

In each of the following examples, assume that the event named `myevent` is defined as shown here:

```
CREATE EVENT myevent
    ON SCHEDULE
      EVERY 6 HOUR
    COMMENT 'A sample comment.'
    DO
      UPDATE myschema.mytable SET mycol = mycol + 1;
```

The following statement changes the schedule for `myevent` from once every six hours starting immediately to once every twelve hours, starting four hours from the time the statement is run:

```
ALTER EVENT myevent
    ON SCHEDULE
      EVERY 12 HOUR
    STARTS CURRENT_TIMESTAMP + INTERVAL 4 HOUR;
```

It is possible to change multiple characteristics of an event in a single statement. This example changes the SQL statement executed by `myevent` to one that deletes all records from `mytable`; it also changes the schedule for the event such that it executes once, one day after this `ALTER EVENT` statement is run.

```
ALTER EVENT myevent
    ON SCHEDULE
      AT CURRENT_TIMESTAMP + INTERVAL 1 DAY
    DO
```

```
      TRUNCATE TABLE myschema.mytable;
```

Specify the options in an `ALTER EVENT` statement only for those characteristics that you want to change; omitted options keep their existing values. This includes any default values for `CREATE EVENT` such as `ENABLE`.

To disable `myevent`, use this `ALTER EVENT` statement:

```
ALTER EVENT myevent
    DISABLE;
```

The `ON SCHEDULE` clause may use expressions involving built-in MySQL functions and user variables to obtain any of the *`timestamp`* or *`interval`* values which it contains. You cannot use stored routines or user-defined functions in such expressions, and you cannot use any table references; however, you can use `SELECT FROM DUAL`. This is true for both `ALTER EVENT` and `CREATE EVENT` statements. References to stored routines, user-defined functions, and tables in such cases are specifically not permitted, and fail with an error (see Bug #22830).

Although an `ALTER EVENT` statement that contains another `ALTER EVENT` statement in its `DO` clause appears to succeed, when the server attempts to execute the resulting scheduled event, the execution fails with an error.

To rename an event, use the `ALTER EVENT` statement's `RENAME TO` clause. This statement renames the event `myevent` to `yourevent`:

```
ALTER EVENT myevent
    RENAME TO yourevent;
```

You can also move an event to a different database using `ALTER EVENT ... RENAME TO ...` and *`db_name.event_name`* notation, as shown here:

```
ALTER EVENT olddb.myevent
    RENAME TO newdb.myevent;
```

To execute the previous statement, the user executing it must have the `EVENT` privilege on both the `olddb` and `newdb` databases.

Note

There is no `RENAME EVENT` statement.

The value `DISABLE ON SLAVE` is used on a replication slave instead of `ENABLE` or `DISABLE` to indicate an event that was created on the master and replicated to the slave, but that is not executed on the slave. Normally, `DISABLE ON SLAVE` is set automatically as required; however, there are some circumstances under which you may want or need to change it manually. See Section 18.4.1.12, “Replication of Invoked Features”, for more information.

13.1.3 ALTER FUNCTION Syntax

```
ALTER FUNCTION func_name [characteristic ...]

characteristic:
    COMMENT 'string'
  | LANGUAGE SQL
  | { CONTAINS SQL | NO SQL | READS SQL DATA | MODIFIES SQL DATA }
  | SQL SECURITY { DEFINER | INVOKER }
```

This statement can be used to change the characteristics of a stored function. More than one change may be specified in an `ALTER FUNCTION` statement. However, you cannot change the parameters or

body of a stored function using this statement; to make such changes, you must drop and re-create the function using `DROP FUNCTION` and `CREATE FUNCTION`.

You must have the `ALTER ROUTINE` privilege for the function. (That privilege is granted automatically to the function creator.) If binary logging is enabled, the `ALTER FUNCTION` statement might also require the `SUPER` privilege, as described in Section 23.7, “Binary Logging of Stored Programs”.

13.1.4 ALTER INSTANCE Syntax

```
ALTER INSTANCE ROTATE INNODB MASTER KEY
```

`ALTER INSTANCE` defines actions applicable to a MySQL server instance.

The `ALTER INSTANCE ROTATE INNODB MASTER KEY` statement is used to rotate the master encryption key used for `InnoDB` tablespace encryption. A keyring plugin must be loaded to use this statement. By default, the MySQL server loads the `keyring_file` plugin. Key rotation requires the `ENCRYPTION_KEY_ADMIN` or `SUPER` privilege.

`ALTER INSTANCE ROTATE INNODB MASTER KEY` supports concurrent DML. However, it cannot be run concurrently with `CREATE TABLE ... ENCRYPTION` or `ALTER TABLE ... ENCRYPTION` operations, and locks are taken to prevent conflicts that could arise from concurrent execution of these statements. If one of the conflicting statements is running, it must complete before another can proceed.

`ALTER INSTANCE` actions are written to the binary log so that they can be executed on replicated servers.

For additional `ALTER INSTANCE ROTATE INNODB MASTER KEY` usage information, see Section 15.7.10, “InnoDB Tablespace Encryption”. For information about the `keyring_file` plugin, see Section 6.5.4, “The MySQL Keyring”.

13.1.5 ALTER PROCEDURE Syntax

```
ALTER PROCEDURE proc_name [characteristic ...]

characteristic:
    COMMENT 'string'
  | LANGUAGE SQL
  | { CONTAINS SQL | NO SQL | READS SQL DATA | MODIFIES SQL DATA }
  | SQL SECURITY { DEFINER | INVOKER }
```

This statement can be used to change the characteristics of a stored procedure. More than one change may be specified in an `ALTER PROCEDURE` statement. However, you cannot change the parameters or body of a stored procedure using this statement; to make such changes, you must drop and re-create the procedure using `DROP PROCEDURE` and `CREATE PROCEDURE`.

You must have the `ALTER ROUTINE` privilege for the procedure. By default, that privilege is granted automatically to the procedure creator. This behavior can be changed by disabling the `automatic_sp_privileges` system variable. See Section 23.2.2, “Stored Routines and MySQL Privileges”.

13.1.6 ALTER SERVER Syntax

```
ALTER SERVER  server_name
    OPTIONS (option [, option] ...)
```

Alters the server information for *`server_name`*, adjusting any of the options permitted in the `CREATE SERVER` statement. The corresponding fields in the `mysql.servers` table are updated accordingly. This statement requires the `SUPER` privilege.

For example, to update the `USER` option:

```
ALTER SERVER s OPTIONS (USER 'sally');
```

`ALTER SERVER` does not cause an automatic commit.

In MySQL 8.0, `ALTER SERVER` is not written to the binary log, regardless of the logging format that is in use.

13.1.7 ALTER TABLE Syntax

```
ALTER TABLE tbl_name
    [alter_specification [, alter_specification] ...]
    [partition_options]

alter_specification:
    table_options
  | ADD [COLUMN] col_name column_definition
        [FIRST | AFTER col_name ]
  | ADD [COLUMN] (col_name column_definition,...)
  | ADD {INDEX|KEY} [index_name]
        [index_type] (index_col_name,...) [index_option] ...
  | ADD [CONSTRAINT [symbol]] PRIMARY KEY
        [index_type] (index_col_name,...) [index_option] ...
  | ADD [CONSTRAINT [symbol]]
        UNIQUE [INDEX|KEY] [index_name]
        [index_type] (index_col_name,...) [index_option] ...
  | ADD FULLTEXT [INDEX|KEY] [index_name]
        (index_col_name,...) [index_option] ...
  | ADD SPATIAL [INDEX|KEY] [index_name]
        (index_col_name,...) [index_option] ...
  | ADD [CONSTRAINT [symbol]]
        FOREIGN KEY [index_name] (index_col_name,...)
        reference_definition
  | ALGORITHM [=] {DEFAULT|INPLACE|COPY}
  | ALTER [COLUMN] col_name {SET DEFAULT literal | DROP DEFAULT}
  | CHANGE [COLUMN] old_col_name new_col_name column_definition
        [FIRST|AFTER col_name]
  | LOCK [=] {DEFAULT|NONE|SHARED|EXCLUSIVE}
  | MODIFY [COLUMN] col_name column_definition
        [FIRST | AFTER col_name]
  | DROP [COLUMN] col_name
  | DROP PRIMARY KEY
  | DROP {INDEX|KEY} index_name
  | DROP FOREIGN KEY fk_symbol
  | ALTER INDEX index_name {VISIBLE | INVISIBLE}
  | DISABLE KEYS
  | ENABLE KEYS
  | RENAME [TO|AS] new_tbl_name
  | RENAME {INDEX|KEY} old_index_name TO new_index_name
  | ORDER BY col_name [, col_name] ...
  | CONVERT TO CHARACTER SET charset_name [COLLATE collation_name]
  | [DEFAULT] CHARACTER SET [=] charset_name [COLLATE [=] collation_name]
  | DISCARD TABLESPACE
  | IMPORT TABLESPACE
  | FORCE
  | {WITHOUT|WITH} VALIDATION
  | ADD PARTITION (partition_definition)
  | DROP PARTITION partition_names
  | DISCARD PARTITION {partition_names | ALL} TABLESPACE
  | IMPORT PARTITION {partition_names | ALL} TABLESPACE
  | TRUNCATE PARTITION {partition_names | ALL}
  | COALESCE PARTITION number
  | REORGANIZE PARTITION partition_names INTO (partition_definitions)
  | EXCHANGE PARTITION partition_name WITH TABLE tbl_name [{WITH|WITHOUT} VALIDATION]
  | ANALYZE PARTITION {partition_names | ALL}
  | CHECK PARTITION {partition_names | ALL}
  | OPTIMIZE PARTITION {partition_names | ALL}
  | REBUILD PARTITION {partition_names | ALL}
  | REPAIR PARTITION {partition_names | ALL}
  | REMOVE PARTITIONING
  | UPGRADE PARTITIONING
```

```
index_col_name:
    col_name [(length)] [ASC | DESC]

index_type:
    USING {BTREE | HASH}

index_option:
    KEY_BLOCK_SIZE [=] value
  | index_type
  | WITH PARSER parser_name
  | COMMENT 'string'
  | {VISIBLE | INVISIBLE}

table_options:
    table_option [[,] table_option] ...

table_option:
    ENGINE [=] engine_name
  | AUTO_INCREMENT [=] value
  | AVG_ROW_LENGTH [=] value
  | [DEFAULT] CHARACTER SET [=] charset_name
  | CHECKSUM [=] {0 | 1}
  | [DEFAULT] COLLATE [=] collation_name
  | COMMENT [=] 'string'
  | COMPRESSION [=] {'ZLIB'|'LZ4'|'NONE'}
  | CONNECTION [=] 'connect_string'
  | DATA DIRECTORY [=] 'absolute path to directory'
  | DELAY_KEY_WRITE [=] {0 | 1}
  | ENCRYPTION [=] {'Y' | 'N'}
  | INDEX DIRECTORY [=] 'absolute path to directory'
  | INSERT_METHOD [=] { NO | FIRST | LAST }
  | KEY_BLOCK_SIZE [=] value
  | MAX_ROWS [=] value
  | MIN_ROWS [=] value
  | PACK_KEYS [=] {0 | 1 | DEFAULT}
  | PASSWORD [=] 'string'
  | ROW_FORMAT [=] {DEFAULT|DYNAMIC|FIXED|COMPRESSED|REDUNDANT|COMPACT}
  | STATS_AUTO_RECALC [=] {DEFAULT|0|1}
  | STATS_PERSISTENT [=] {DEFAULT|0|1}
  | STATS_SAMPLE_PAGES [=] value
  | TABLESPACE tablespace_name
  | UNION [=] (tbl_name[,tbl_name]...)

partition_options:
    (see CREATE TABLE options)
```

`ALTER TABLE` changes the structure of a table. For example, you can add or delete columns, create or destroy indexes, change the type of existing columns, or rename columns or the table itself. You can also change characteristics such as the storage engine used for the table or the table comment.

- To use `ALTER TABLE`, you need `ALTER`, `CREATE`, and `INSERT` privileges for the table. Renaming a table requires `ALTER` and `DROP` on the old table, `ALTER`, `CREATE`, and `INSERT` on the new table.
- Following the table name, specify the alterations to be made. If none are given, `ALTER TABLE` does nothing.
- The syntax for many of the permissible alterations is similar to clauses of the `CREATE TABLE` statement. *`column_definition`* clauses use the same syntax for `ADD` and `CHANGE` as for `CREATE TABLE`. See Section 13.1.15, “CREATE TABLE Syntax”, for more information.
- The word `COLUMN` is optional and can be omitted.
- You can issue multiple `ADD`, `ALTER`, `DROP`, and `CHANGE` clauses in a single `ALTER TABLE` statement, separated by commas. This is a MySQL extension to standard SQL, which permits only one of each clause per `ALTER TABLE` statement. For example, to drop multiple columns in a single statement, do this:

```
ALTER TABLE t2 DROP COLUMN c, DROP COLUMN d;
```

- Some operations may result in warnings if attempted on a table for which the storage engine does not support the operation. These warnings can be displayed with `SHOW WARNINGS`. See Section 13.7.5.40, “SHOW WARNINGS Syntax”. For information on troubleshooting `ALTER TABLE`, see Section B.5.6.1, “Problems with ALTER TABLE”.
- For usage examples, see Section 13.1.7.3, “ALTER TABLE Examples”.
- For information about generated columns, see Section 13.1.7.2, “ALTER TABLE and Generated Columns”.
- With the `mysql_info()` C API function, you can find out how many rows were copied by `ALTER TABLE`. See Section 27.7.7.36, “mysql_info()”.

There are several additional aspects to the `ALTER TABLE` statement, described under the following topics in this section:

- Table Options
- Performance and Storage Considerations
- Using the LOCK Clause to Control Concurrency
- Column Attributes
- Primary Keys and Indexes
- Foreign Keys
- Dropping Columns
- Changing the Character Set
- Discarding and Importing InnoDB Tablespaces
- Row Order for MyISAM Tables
- Partitioning Options

Table Options

`table_options` signifies table options of the kind that can be used in the `CREATE TABLE` statement, such as `ENGINE`, `AUTO_INCREMENT`, `AVG_ROW_LENGTH`, `MAX_ROWS`, `ROW_FORMAT`, or `TABLESPACE`.

For descriptions of all table options, see Section 13.1.15, “CREATE TABLE Syntax”. However, `ALTER TABLE` ignores `DATA DIRECTORY` and `INDEX DIRECTORY` when given as table options. `ALTER TABLE` permits them only as partitioning options, and requires that you have the `FILE` privilege.

Use of table options with `ALTER TABLE` provides a convenient way of altering single table characteristics. For example:

- If `t1` is currently not an `InnoDB` table, this statement changes its storage engine to `InnoDB`:

```
ALTER TABLE t1 ENGINE = InnoDB;
```

  - See Section 15.8.1.4, “Converting Tables from MyISAM to InnoDB” for considerations when switching tables to the `InnoDB` storage engine.
  - When you specify an `ENGINE` clause, `ALTER TABLE` rebuilds the table. This is true even if the table already has the specified storage engine.
  - Running `ALTER TABLE *tbl_name* ENGINE=INNODB` on an existing `InnoDB` table performs a “null” `ALTER TABLE` operation, which can be used to defragment an `InnoDB` table, as described

in Section 15.11.4, “Defragmenting a Table”. Running `ALTER TABLE tbl_name FORCE` on an `InnoDB` table performs the same function.

- `ALTER TABLE tbl_name ENGINE=INNODB` and `ALTER TABLE tbl_name FORCE` use online DDL. For more information, see Section 15.12.1, “Online DDL Overview”.
- The outcome of attempting to change the storage engine of a table is affected by whether the desired storage engine is available and the setting of the `NO_ENGINE_SUBSTITUTION` SQL mode, as described in Section 5.1.8, “Server SQL Modes”.
- To prevent inadvertent loss of data, `ALTER TABLE` cannot be used to change the storage engine of a table to `MERGE` or `BLACKHOLE`.

- To change the `InnoDB` table to use compressed row-storage format:

```
ALTER TABLE t1 ROW_FORMAT = COMPRESSED;
```

- If the `InnoDB` tablespace encryption feature is enabled (see Section 15.7.10, “InnoDB Tablespace Encryption”), encryption for `t1` can be enabled or disabled like this:

```
ALTER TABLE t1 ENCRYPTION='Y';
ALTER TABLE t1 ENCRYPTION='N';
```

- To reset the current auto-increment value:

```
ALTER TABLE t1 AUTO_INCREMENT = 13;
```

  You cannot reset the counter to a value less than or equal to the value that is currently in use. For both `InnoDB` and `MyISAM`, if the value is less than or equal to the maximum value currently in the `AUTO_INCREMENT` column, the value is reset to the current maximum `AUTO_INCREMENT` column value plus one.

- To change the default table character set:

```
ALTER TABLE t1 CHARACTER SET = utf8;
```

- To add (or change) a table comment:

```
ALTER TABLE t1 COMMENT = 'New table comment';
```

- You can use `ALTER TABLE` with the `TABLESPACE` option to move non-partitioned `InnoDB` tables between existing general tablespaces, file-per-table tablespaces, and the system tablespace. See Moving Non-Partitioned Tables Between Tablespaces Using ALTER TABLE.
  - For partitioned tables, `ALTER TABLE tbl_name TABLESPACE [=] tablespace_name` only modifies the default tablespace. It does not move partitions from one tablespace to another. To move table partitions, you must move each partition using `ALTER TABLE tbl_name REORGANIZE PARTITION`. See Moving Table Partitions Between Tablespaces Using ALTER TABLE.
  - `ALTER TABLE ... TABLESPACE` operations always cause a full table rebuild, even if the `TABLESPACE` attribute has not changed from its previous value.
  - `ALTER TABLE ... TABLESPACE` syntax does not support moving a table from a temporary tablespace to a persistent tablespace.
  - The `DATA DIRECTORY` clause, which is supported with `CREATE TABLE ... TABLESPACE`, is not supported with `ALTER TABLE ... TABLESPACE`, and is ignored if specified.

- For more information about the capabilities and limitations of the `TABLESPACE` option, see `CREATE TABLE`.

To verify that the table options were changed as intended, use `SHOW CREATE TABLE`, or query `INFORMATION_SCHEMA.TABLES`.

Performance and Storage Considerations

`ALTER TABLE` operations that are **not** performed in place make a temporary copy of the original table. MySQL waits for other operations that are modifying the table, then proceeds. It incorporates the alteration into the copy, deletes the original table, and renames the new one. While `ALTER TABLE` is executing, the original table is readable by other sessions (with the exception noted shortly). Updates and writes to the table that begin after the `ALTER TABLE` operation begins are stalled until the new table is ready, then are automatically redirected to the new table without any failed updates. The temporary copy of the original table is created in the database directory of the new table. This can differ from the database directory of the original table for `ALTER TABLE` operations that rename the table to a different database.

The exception referred to earlier is that `ALTER TABLE` blocks reads (not just writes) at the point where it is ready to clear outdated table structures from the table and table definition caches. At this point, it must acquire an exclusive lock. To do so, it waits for current readers to finish, and blocks new reads (and writes).

For `MyISAM` tables, you can speed up index re-creation (the slowest part of the alteration process) by setting the `myisam_sort_buffer_size` system variable to a high value.

For `InnoDB` tables, a table-copying `ALTER TABLE` operation on table that resides in a shared tablespace such as a general tablespace or the system tablespace can increase the amount of space used by the tablespace. Such operations require as much additional space as the data in the table plus indexes. For a table that resides in a shared tablespace, the additional space used during a table-copying `ALTER TABLE` operation is not released back to the operating system as it is for a table that resides in a file-per-table tablespace.

`ALTER TABLE` operations that are performed in place do not require creating a temporary copy of the original table. These operations include:

- `ALTER TABLE` operations on `InnoDB` tables that are supported by the `InnoDB` online DDL feature. For an overview of supported operations, see Section 15.12.1, “Online DDL Overview”. For information about performance and concurrency of online DDL operations, see Section 15.12.2, “Online DDL Performance, Concurrency, and Space Requirements”.

- `ALTER TABLE` *`tbl_name`* `RENAME TO` *`new_tbl_name`*. When run without other options, MySQL renames files that correspond to the table *`tbl_name`* without making a copy. (You can also use the `RENAME TABLE` statement to rename tables. See Section 13.1.29, “RENAME TABLE Syntax”.) Privileges granted specifically for the renamed table are not migrated to the new name. They must be changed manually.

- Alterations that modify only table metadata and not table data are immediate because the server only needs to alter table metadata, not touch table contents. The following changes are made in this way:

  - Renaming a column.

  - Changing the default value of a column.

  - Changing the definition of an `ENUM` or `SET` column by adding new enumeration or set members to the *end* of the list of valid member values, as long as the storage size of the data type does not change. For example, adding a member to a `SET` column that has 8 members changes the required storage per value from 1 byte to 2 bytes; this will require a table copy. Adding members in the middle of the list causes renumbering of existing members, which requires a table copy.

- Renaming an index.
- Adding or dropping an index, for InnoDB. See Section 15.12.1, “Online DDL Overview”.
- Modifying index visibility with an ALTER INDEX operation.
- Column modifications of tables containing generated columns that depend on columns with a DEFAULT value if the modified columns are not involved in the generated column expressions. For example, changing the NULL property of a separate column can be done in place without a table rebuild.

Specifying ALGORITHM=INPLACE makes the operation use the in-place technique for clauses and storage engines that support it, and fail with an error otherwise, thus avoiding a lengthy table copy if you try altering a table that uses a different storage engine than you expect.

You can force an ALTER TABLE operation that would otherwise not use the table copy method by setting the old_alter_table system variable to ON, or specifying ALGORITHM=COPY as one of the *alter_specification* clauses. If there is a conflict between the old_alter_table setting and an ALGORITHM clause with a value other than DEFAULT, the ALGORITHM clause takes precedence.

Specifying ALGORITHM=DEFAULT is the same a specifying no ALGORITHM clause at all, in which case ALGORITM=INPLACE is used if supported by the storage engine. Otherwise, ALGORITHM=COPY is used.

An ALTER TABLE operation run with the ALGORITHM=COPY clause prevents concurrent DML operations. Concurrent queries are still allowed. That is, a table-copying operation always includes at least the concurrency restrictions of LOCK=SHARED (allow queries but not DML). You can further restrict concurrency for such operations by specifying LOCK=EXCLUSIVE, which prevents DML and queries.

As of MySQL 5.7.4, ALTER TABLE upgrades MySQL 5.5 temporal columns to 5.6 format for ADD COLUMN, CHANGE COLUMN, MODIFY COLUMN, ADD INDEX, and FORCE operations. This conversion cannot be done using the INPLACE algorithm because the table must be rebuilt, so specifying ALGORITHM=INPLACE in these cases results in an error. Specify ALGORITHM=COPY if necessary.

If an ALTER TABLE operation on a multicolumn index used to partition a table by KEY changes the order of the columns, it can only be performed using ALGORITHM=COPY.

The WITHOUT VALIDATION and WITH VALIDATION clauses affect whether ALTER TABLE performs an in-place operation for virtual generated column modifications. See Section 13.1.7.2, “ALTER TABLE and Generated Columns”.

ALTER TABLE with DISCARD ... PARTITION ... TABLESPACE or IMPORT ... PARTITION ... TABLESPACE does not create any temporary tables or temporary partition files.

ALTER TABLE with ADD PARTITION, DROP PARTITION, COALESCE PARTITION, REBUILD PARTITION, or REORGANIZE PARTITION does not create temporary tables (except when used with NDB tables); however, these operations can and do create temporary partition files.

ADD or DROP operations for RANGE or LIST partitions are immediate operations or nearly so. ADD or COALESCE operations for HASH or KEY partitions copy data between all partitions, unless LINEAR HASH or LINEAR KEY was used; this is effectively the same as creating a new table, although the ADD or COALESCE operation is performed partition by partition. REORGANIZE operations copy only changed partitions and do not touch unchanged ones.

Using the LOCK Clause to Control Concurrency

You can control the level of concurrent reading and writing of the table while it is being altered, using the LOCK clause. Specifying a non-default value for this clause lets you require a certain amount of

concurrent access or exclusivity during the alter operation, and halts the operation if the requested degree of locking is not available. The parameters for the `LOCK` clause are:

- `LOCK = DEFAULT`

  Maximum level of concurrency for the given `ALGORITHM` clause (if any) and `ALTER TABLE` operation: Permit concurrent reads and writes if supported. If not, permit concurrent reads if supported. If not, enforce exclusive access.

- `LOCK = NONE`

  If supported, permit concurrent reads and writes. Otherwise, return an error message.

- `LOCK = SHARED`

  If supported, permit concurrent reads but block writes. Note that writes will be blocked even if concurrent writes are supported by the storage engine for the given `ALGORITHM` clause (if any) and `ALTER TABLE` operation. If concurrent reads are not supported, return an error message.

- `LOCK = EXCLUSIVE`

  Enforce exclusive access. This will be done even if concurrent reads/writes are supported by the storage engine for the given `ALGORITHM` clause (if any) and `ALTER TABLE` operation.

Column Attributes

- You can rename a column using a `CHANGE old_col_name new_col_name column_definition` clause. To do so, specify the old and new column names and the definition that the column currently has. For example, to rename an `INTEGER` column from `a` to `b`, you can do this:

  ```
  ALTER TABLE t1 CHANGE a b INTEGER;
  ```

- To change a column's type but not the name, `CHANGE` syntax still requires an old and new column name, even if they are the same. For example:

  ```
  ALTER TABLE t1 CHANGE b b BIGINT NOT NULL;
  ```

- You can also use `MODIFY` to change a column's type without renaming it:

  ```
  ALTER TABLE t1 MODIFY b BIGINT NOT NULL;
  ```

  `MODIFY` is an extension to `ALTER TABLE` for Oracle compatibility.

- When you use `CHANGE` or `MODIFY`, `column_definition` must include the data type and all attributes that should apply to the new column, other than index attributes such as `PRIMARY KEY` or `UNIQUE`. Attributes present in the original definition but not specified for the new definition are not carried forward. Suppose that a column `col1` is defined as `INT UNSIGNED DEFAULT 1 COMMENT 'my column'` and you modify the column as follows:

  ```
  ALTER TABLE t1 MODIFY col1 BIGINT;
  ```

  The resulting column will be defined as `BIGINT`, but will not include the attributes `UNSIGNED DEFAULT 1 COMMENT 'my column'`. To retain them, the statement should be:

  ```
  ```

```
ALTER TABLE t1 MODIFY col1 BIGINT UNSIGNED DEFAULT 1 COMMENT 'my column';
```

- When you change a data type using `CHANGE` or `MODIFY`, MySQL tries to convert existing column values to the new type as well as possible.

Warning

This conversion may result in alteration of data. For example, if you shorten a string column, values may be truncated. To prevent the operation from succeeding if conversions to the new data type would result in loss of data, enable strict SQL mode before using `ALTER TABLE` (see Section 5.1.8, “Server SQL Modes”).

- To add a column at a specific position within a table row, use `FIRST` or `AFTER col_name`. The default is to add the column last. You can also use `FIRST` and `AFTER` in `CHANGE` or `MODIFY` operations to reorder columns within a table.

- If you use `CHANGE` or `MODIFY` to shorten a column for which an index exists on the column, and the resulting column length is less than the index length, MySQL shortens the index automatically.

- `CHANGE col_name` is a MySQL extension to standard SQL.

- `ALTER ... SET DEFAULT` or `ALTER ... DROP DEFAULT` specify a new default value for a column or remove the old default value, respectively. If the old default is removed and the column can be `NULL`, the new default is `NULL`. If the column cannot be `NULL`, MySQL assigns a default value as described in Section 11.7, “Data Type Default Values”.

Primary Keys and Indexes

- `DROP PRIMARY KEY` drops the primary key. If there is no primary key, an error occurs. For information about the performance characteristics of primary keys, especially for `InnoDB` tables, see Section 8.3.2, “Using Primary Keys”.

- If you add a `UNIQUE INDEX` or `PRIMARY KEY` to a table, MySQL stores it before any nonunique index to permit detection of duplicate keys as early as possible.

- `DROP INDEX` removes an index. This is a MySQL extension to standard SQL. See Section 13.1.22, “DROP INDEX Syntax”. If you are unsure of the index name, use `SHOW INDEX FROM tbl_name`.

- Some storage engines permit you to specify an index type when creating an index. The syntax for the *`index_type`* specifier is `USING type_name`. For details about `USING`, see Section 13.1.12, “CREATE INDEX Syntax”. The preferred position is after the column list. Support for use of the option before the column list will be removed in a future MySQL release.

- *`index_option`* values specify additional options for an index. `USING` is one such option. For details about permissible *`index_option`* values, see Section 13.1.12, “CREATE INDEX Syntax”.

- `RENAME INDEX old_index_name TO new_index_name` renames an index. This is a MySQL extension to standard SQL. The content of the table remains unchanged. *`old_index_name`* must be the name of an existing index in the table that is not dropped by the same `ALTER TABLE` statement. `new_index_name` is the new index name, which cannot duplicate the name of an index in the resulting table after changes have been applied. Neither index name can be `PRIMARY`.

- If you use `ALTER TABLE` on a `MyISAM` table, all nonunique indexes are created in a separate batch (as for `REPAIR TABLE`). This should make `ALTER TABLE` much faster when you have many indexes.

- For `MyISAM` tables, key updating can be controlled explicitly. Use `ALTER TABLE ... DISABLE KEYS` to tell MySQL to stop updating nonunique indexes. Then use `ALTER TABLE ... ENABLE KEYS` to re-create missing indexes. `MyISAM` does this with a special algorithm that is much faster

than inserting keys one by one, so disabling keys before performing bulk insert operations should give a considerable speedup. Using `ALTER TABLE ... DISABLE KEYS` requires the `INDEX` privilege in addition to the privileges mentioned earlier.

While the nonunique indexes are disabled, they are ignored for statements such as `SELECT` and `EXPLAIN` that otherwise would use them.

- After an `ALTER TABLE` statement, it may be necessary to run `ANALYZE TABLE` to update index cardinality information. See Section 13.7.5.22, “SHOW INDEX Syntax”.
- The `ALTER INDEX` operation permits an index to be made visible or invisible. An invisible index is not used by the optimizer. Modification of index visibility applies to indexes other than primary keys (either explicit or implicit). This feature is storage engine neutral (supported for any engine). For more information, see Section 8.3.11, “Invisible Indexes”.

Foreign Keys

- The `FOREIGN KEY` and `REFERENCES` clauses are supported by the `InnoDB` storage engine, which implements `ADD [CONSTRAINT [symbol]] FOREIGN KEY [index_name] (...) REFERENCES ... (...)`. See Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”. For other storage engines, the clauses are parsed but ignored. The `CHECK` clause is parsed but ignored by all storage engines. See Section 13.1.15, “CREATE TABLE Syntax”. The reason for accepting but ignoring syntax clauses is for compatibility, to make it easier to port code from other SQL servers, and to run applications that create tables with references. See Section 1.8.2, “MySQL Differences from Standard SQL”.

  For `ALTER TABLE`, unlike `CREATE TABLE`, `ADD FOREIGN KEY` ignores *`index_name`* if given and uses an automatically generated foreign key name. As a workaround, include the `CONSTRAINT` clause to specify the foreign key name:

```
ADD CONSTRAINT name FOREIGN KEY (....) ...
```

Important

The inline `REFERENCES` specifications where the references are defined as part of the column specification are silently ignored. MySQL only accepts `REFERENCES` clauses defined as part of a separate `FOREIGN KEY` specification.

Note

Partitioned `InnoDB` tables do not support foreign keys. See Section 22.6.2, “Partitioning Limitations Relating to Storage Engines”, for more information.

- MySQL supports the use of `ALTER TABLE` to drop foreign keys:

```
ALTER TABLE tbl_name DROP FOREIGN KEY fk_symbol;
```

- Adding and dropping a foreign key in the same `ALTER TABLE` statement is supported for `ALTER TABLE ... ALGORITHM=INPLACE` but is unsupported for `ALTER TABLE ... ALGORITHM=COPY`.
- The server prohibits changes to foreign key columns that have the potential to cause loss of referential integrity. It also prohibits changes to the data type of such columns that may be unsafe. For example, changing `VARCHAR(20)` to `VARCHAR(30)` is permitted, but changing it to `VARCHAR(1024)` is not because that alters the number of length bytes required to store individual values. A workaround is to use `ALTER TABLE ... DROP FOREIGN KEY` before changing the column definition and `ALTER TABLE ... ADD FOREIGN KEY` afterward.
- `ALTER TABLE tbl_name RENAME new_tbl_name` changes internally generated foreign key constraint names and user-defined foreign key constraint names that contain the string

"*tbl_name*_ibfk_" to reflect the new table name. InnoDB interprets foreign key constraint names that contain the string "*tbl_name*_ibfk_" as internally generated names.

Dropping Columns

- If a table contains only one column, the column cannot be dropped. If what you intend is to remove the table, use DROP TABLE instead.
- If columns are dropped from a table, the columns are also removed from any index of which they are a part. If all columns that make up an index are dropped, the index is dropped as well. If you use CHANGE or MODIFY to shorten a column for which an index exists on the column, and the resulting column length is less than the index length, MySQL shortens the index automatically.
- DROP *col_name* is a MySQL extension to standard SQL.

Changing the Character Set

To change the table default character set and all character columns (CHAR, VARCHAR, TEXT) to a new character set, use a statement like this:

```
ALTER TABLE tbl_name CONVERT TO CHARACTER SET charset_name;
```

The statement also changes the collation of all character columns. If you specify no COLLATE clause to indicate which collation to use, the statement uses default collation for the character set. If this collation is inappropriate for the intended table use (for example, if it would change from a case-sensitive collation to a case-insensitive collation), specify a collation explicitly.

For a column that has a data type of VARCHAR or one of the TEXT types, CONVERT TO CHARACTER SET will change the data type as necessary to ensure that the new column is long enough to store as many characters as the original column. For example, a TEXT column has two length bytes, which store the byte-length of values in the column, up to a maximum of 65,535. For a latin1 TEXT column, each character requires a single byte, so the column can store up to 65,535 characters. If the column is converted to utf8, each character might require up to three bytes, for a maximum possible length of 3 × 65,535 = 196,605 bytes. That length will not fit in a TEXT column's length bytes, so MySQL will convert the data type to MEDIUMTEXT, which is the smallest string type for which the length bytes can record a value of 196,605. Similarly, a VARCHAR column might be converted to MEDIUMTEXT.

To avoid data type changes of the type just described, do not use CONVERT TO CHARACTER SET. Instead, use MODIFY to change individual columns. For example:

```
ALTER TABLE t MODIFY latin1_text_col TEXT CHARACTER SET utf8;
ALTER TABLE t MODIFY latin1_varchar_col VARCHAR(M) CHARACTER SET utf8;
```

If you specify CONVERT TO CHARACTER SET binary, the CHAR, VARCHAR, and TEXT columns are converted to their corresponding binary string types (BINARY, VARBINARY, BLOB). This means that the columns no longer will have a character set and a subsequent CONVERT TO operation will not apply to them.

If *charset_name* is DEFAULT, the database character set is used.

Warning

The CONVERT TO operation converts column values between the character sets. This is *not* what you want if you have a column in one character set (like latin1) but the stored values actually use some other, incompatible character set (like utf8). In this case, you have to do the following for each such column:

```
ALTER TABLE t1 CHANGE c1 c1 BLOB;
ALTER TABLE t1 CHANGE c1 c1 TEXT CHARACTER SET utf8;
```

> The reason this works is that there is no conversion when you convert to or from `BLOB` columns.

To change only the *default* character set for a table, use this statement:

```
ALTER TABLE tbl_name DEFAULT CHARACTER SET charset_name;
```

The word `DEFAULT` is optional. The default character set is the character set that is used if you do not specify the character set for columns that you add to a table later (for example, with `ALTER TABLE ... ADD column`).

When `foreign_key_checks` is enabled, which is the default setting, character set conversion is not permitted on tables that include a character string column used in a foreign key constraint. The workaround is to disable `foreign_key_checks` before performing the character set conversion. You must perform the conversion on both tables involved in the foreign key constraint before re-enabling `foreign_key_checks`. If you re-enable `foreign_key_checks` after converting only one of the tables, an `ON DELETE CASCADE` or `ON UPDATE CASCADE` operation could corrupt data in the referencing table due to implicit conversion that occurs during these operations (Bug #45290, Bug #74816).

Discarding and Importing InnoDB Tablespaces

An `InnoDB` table created in its own file-per-table tablespace can be discarded and imported using the `DISCARD TABLESPACE` and `IMPORT TABLESPACE` options. These options can be used to import a file-per-table tablespace from a backup or to copy a file-per-table tablespace from one database server to another. See Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance”.

Row Order for MyISAM Tables

`ORDER BY` enables you to create the new table with the rows in a specific order. This option is useful primarily when you know that you query the rows in a certain order most of the time. By using this option after major changes to the table, you might be able to get higher performance. In some cases, it might make sorting easier for MySQL if the table is in order by the column that you want to order it by later.

> **Note**
>
> The table does not remain in the specified order after inserts and deletes.

`ORDER BY` syntax permits one or more column names to be specified for sorting, each of which optionally can be followed by `ASC` or `DESC` to indicate ascending or descending sort order, respectively. The default is ascending order. Only column names are permitted as sort criteria; arbitrary expressions are not permitted. This clause should be given last after any other clauses.

`ORDER BY` does not make sense for `InnoDB` tables because `InnoDB` always orders table rows according to the clustered index.

When used on a partitioned table, `ALTER TABLE ... ORDER BY` orders rows within each partition only.

Partitioning Options

`partition_options` signifies options that can be used with partitioned tables for repartitioning, for adding, dropping, discarding, importing, merging, and splitting partitions, and for performing partitioning maintenance.

It is possible for an `ALTER TABLE` statement to contain a `PARTITION BY` or `REMOVE PARTITIONING` clause in an addition to other alter specifications, but the `PARTITION BY` or `REMOVE PARTITIONING` clause must be specified last after any other specifications. The `ADD PARTITION`, `DROP PARTITION`,

DISCARD PARTITION, IMPORT PARTITION, COALESCE PARTITION, REORGANIZE PARTITION, EXCHANGE PARTITION, ANALYZE PARTITION, CHECK PARTITION, and REPAIR PARTITION options cannot be combined with other alter specifications in a single ALTER TABLE, since the options just listed act on individual partitions.

For more information about partition options, see Section 13.1.15, “CREATE TABLE Syntax”, and Section 13.1.7.1, “ALTER TABLE Partition Operations”. For information about and examples of ALTER TABLE ... EXCHANGE PARTITION statements, see Section 22.3.3, “Exchanging Partitions and Subpartitions with Tables”.

Prior to MySQL 5.7.6, partitioned InnoDB tables used the generic ha_partition partitioning handler employed by MyISAM and other storage engines not supplying their own partitioning handlers; in MySQL 5.7.6 and later, such tables are created using the InnoDB storage engine's own (or “native”) partitioning handler. Beginning with MySQL 5.7.9, you can upgrade an InnoDB table that was created in MySQL 5.7.6 or earlier (that is, created using ha_partition) to the InnoDB native partition handler using ALTER TABLE ... UPGRADE PARTITIONING. (Bug #76734, Bug #20727344) This version of ALTER TABLE does not accept any other options and can be used only on a single table at a time. You can also use mysql_upgrade to upgrade older partitioned InnoDB tables to the native partitioning handler.

13.1.7.1 ALTER TABLE Partition Operations

Partitioning-related clauses for ALTER TABLE can be used with partitioned tables for repartitioning, for adding, dropping, discarding, importing, merging, and splitting partitions, and for performing partitioning maintenance.

- Simply using a *partition_options* clause with ALTER TABLE on a partitioned table repartitions the table according to the partitioning scheme defined by the *partition_options*. This clause always begins with PARTITION BY, and follows the same syntax and other rules as apply to the *partition_options* clause for CREATE TABLE (see Section 13.1.15, “CREATE TABLE Syntax”, for more detailed information), and can also be used to partition an existing table that is not already partitioned. For example, consider a (nonpartitioned) table defined as shown here:

```
CREATE TABLE t1 (
    id INT,
    year_col INT
);
```

This table can be partitioned by HASH, using the id column as the partitioning key, into 8 partitions by means of this statement:

```
ALTER TABLE t1
    PARTITION BY HASH(id)
    PARTITIONS 8;
```

MySQL supports an ALGORITHM option with [SUB]PARTITION BY [LINEAR] KEY. ALGORITHM=1 causes the server to use the same key-hashing functions as MySQL 5.1 when computing the placement of rows in partitions; ALGORITHM=2 means that the server employs the key-hashing functions implemented and used by default for new KEY partitioned tables in MySQL 5.5 and later. (Partitioned tables created with the key-hashing functions employed in MySQL 5.5 and later cannot be used by a MySQL 5.1 server.) Not specifying the option has the same effect as using ALGORITHM=2. This option is intended for use chiefly when upgrading or downgrading [LINEAR] KEY partitioned tables between MySQL 5.1 and later MySQL versions, or for creating tables partitioned by KEY or LINEAR KEY on a MySQL 5.5 or later server which can be used on a MySQL 5.1 server.

To upgrade a KEY partitioned table that was created in MySQL 5.1, first execute SHOW CREATE TABLE and note the exact columns and number of partitions shown. Now execute an ALTER TABLE statement using exactly the same column list and number of partitions as in the CREATE TABLE statement, while adding ALGORITHM=2 immediately following the PARTITION BY keywords. (You

should also include the `LINEAR` keyword if it was used for the original table definition.) An example from a session in the `mysql` client is shown here:

```
mysql> SHOW CREATE TABLE p\G
*************************** 1. row ***************************
       Table: p
Create Table: CREATE TABLE `p` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `cd` datetime NOT NULL,
  PRIMARY KEY (`id`)
) ENGINE=InnoDB DEFAULT CHARSET=latin1
/*!50100 PARTITION BY LINEAR KEY (id)
PARTITIONS 32 */
1 row in set (0.00 sec)

mysql> ALTER TABLE p PARTITION BY LINEAR KEY ALGORITHM=2 (id) PARTITIONS 32;
Query OK, 0 rows affected (5.34 sec)
Records: 0  Duplicates: 0  Warnings: 0

mysql> SHOW CREATE TABLE p\G
*************************** 1. row ***************************
       Table: p
Create Table: CREATE TABLE `p` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `cd` datetime NOT NULL,
  PRIMARY KEY (`id`)
) ENGINE=InnoDB DEFAULT CHARSET=latin1
/*!50100 PARTITION BY LINEAR KEY (id)
PARTITIONS 32 */
1 row in set (0.00 sec)
```

Downgrading a table created using the default key-hashing used in MySQL 5.5 and later to enable its use by a MySQL 5.1 server is similar, except in this case you should use `ALGORITHM=1` to force the table's partitions to be rebuilt using the MySQL 5.1 key-hashing functions. It is recommended that you not do this except when necessary for compatibility with a MySQL 5.1 server, as the improved `KEY` hashing functions used by default in MySQL 5.5 and later provide fixes for a number of issues found in the older implementation.

Note

A table upgraded by means of `ALTER TABLE ... PARTITION BY ALGORITHM=2 [LINEAR] KEY ...` can no longer be used by a MySQL 5.1 server. (Such a table would need to be downgraded with `ALTER TABLE ... PARTITION BY ALGORITHM=1 [LINEAR] KEY ...` before it could be used again by a MySQL 5.1 server.)

The table that results from using an `ALTER TABLE ... PARTITION BY` statement must follow the same rules as one created using `CREATE TABLE ... PARTITION BY`. This includes the rules governing the relationship between any unique keys (including any primary key) that the table might have, and the column or columns used in the partitioning expression, as discussed in Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”. The `CREATE TABLE ... PARTITION BY` rules for specifying the number of partitions also apply to `ALTER TABLE ... PARTITION BY`.

The *`partition_definition`* clause for `ALTER TABLE ADD PARTITION` supports the same options as the clause of the same name for the `CREATE TABLE` statement. (See Section 13.1.15, “CREATE TABLE Syntax”, for the syntax and description.) Suppose that you have the partitioned table created as shown here:

```
CREATE TABLE t1 (
    id INT,
    year_col INT
)
PARTITION BY RANGE (year_col) (
    PARTITION p0 VALUES LESS THAN (1991),
```

```
    PARTITION p1 VALUES LESS THAN (1995),
    PARTITION p2 VALUES LESS THAN (1999)
);
```

You can add a new partition `p3` to this table for storing values less than `2002` as follows:

```
ALTER TABLE t1 ADD PARTITION (PARTITION p3 VALUES LESS THAN (2002));
```

`ADD PARTITION` can also be used with the `TABLESPACE` clause to add a new partition to an existing general tablespace, to a file-per-table tablespace, or to the system tablespace.

```
ALTER TABLE t1 ADD PARTITION (PARTITION p4 VALUES LESS THAN (2015) TABLESPACE = `ts1`);

ALTER TABLE t1 ADD PARTITION (PARTITION p4 VALUES LESS THAN (2015) TABLESPACE = `innodb_file_per_tabl

ALTER TABLE t1 ADD PARTITION (PARTITION p4 VALUES LESS THAN (2015) TABLESPACE = `innodb_system`);
```

Note

If the `TABLESPACE = tablespace_name` option is not defined, the `ALTER TABLE ... ADD PARTITION` operation adds the partition to the table's default tablespace, which can be specified at the table level during `CREATE TABLE` or `ALTER TABLE`.

`DROP PARTITION` can be used to drop one or more `RANGE` or `LIST` partitions. This statement cannot be used with `HASH` or `KEY` partitions; instead, use `COALESCE PARTITION` (see later in this section). Any data that was stored in the dropped partitions named in the *`partition_names`* list is discarded. For example, given the table `t1` defined previously, you can drop the partitions named `p0` and `p1` as shown here:

```
ALTER TABLE t1 DROP PARTITION p0, p1;
```

`ADD PARTITION` and `DROP PARTITION` do not currently support `IF [NOT] EXISTS`.

The `DISCARD PARTITION ... TABLESPACE` and `IMPORT PARTITION ... TABLESPACE` options extend the Transportable Tablespace feature to individual `InnoDB` table partitions. Each `InnoDB` table partition has its own tablespace file (`.idb` file). The Transportable Tablespace feature makes it easy to copy the tablespaces from a running MySQL server instance to another running instance, or to perform a restore on the same instance. Both options take a comma-separated list of one or more partition names. For example:

```
ALTER TABLE t1 DISCARD PARTITION p2, p3 TABLESPACE;

ALTER TABLE t1 IMPORT PARTITION p2, p3 TABLESPACE;
```

When running `DISCARD PARTITION ... TABLESPACE` and `IMPORT PARTITION ... TABLESPACE` on subpartitioned tables, both partition and subpartition names are allowed. When a partition name is specified, subpartitions of that partition are included.

The Transportable Tablespace feature also supports copying or restoring partitioned `InnoDB` tables (all partitions at once). For additional information, see Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance”, as well as, Section 15.7.6.1, “Transportable Tablespace Examples”.

Renames of partitioned table are supported. You can rename individual partitions indirectly using `ALTER TABLE ... REORGANIZE PARTITION`; however, this operation copies the partition's data.

It is possible to delete rows from selected partitions using the `TRUNCATE PARTITION` option. This option takes a comma-separated list of one or more partition names. Consider the table `t1` created by this statement:

```
CREATE TABLE t1 (
    id INT,
    year_col INT
)
PARTITION BY RANGE (year_col) (
    PARTITION p0 VALUES LESS THAN (1991),
    PARTITION p1 VALUES LESS THAN (1995),
    PARTITION p2 VALUES LESS THAN (1999),
    PARTITION p3 VALUES LESS THAN (2003),
    PARTITION p4 VALUES LESS THAN (2007)
);
```

To delete all rows from partition `p0`, you can use the following statement:

```
ALTER TABLE t1 TRUNCATE PARTITION p0;
```

The statement just shown has the same effect as the following `DELETE` statement:

```
DELETE FROM t1 WHERE year_col < 1991;
```

When truncating multiple partitions, the partitions do not have to be contiguous: This can greatly simplify delete operations on partitioned tables that would otherwise require very complex `WHERE` conditions if done with `DELETE` statements. For example, this statement deletes all rows from partitions `p1` and `p3`:

```
ALTER TABLE t1 TRUNCATE PARTITION p1, p3;
```

An equivalent `DELETE` statement is shown here:

```
DELETE FROM t1 WHERE
    (year_col >= 1991 AND year_col < 1995)
    OR
    (year_col >= 2003 AND year_col < 2007);
```

You can use the `ALL` keyword in place of the list of partition names; in this case, the statement acts on all partitions in the table.

`TRUNCATE PARTITION` merely deletes rows; it does not alter the definition of the table itself, or of any of its partitions.

You can verify that the rows were dropped by checking the `INFORMATION_SCHEMA.PARTITIONS` table, using a query such as this one:

```
SELECT PARTITION_NAME, TABLE_ROWS
    FROM INFORMATION_SCHEMA.PARTITIONS
    WHERE TABLE_NAME = 't1';
```

`COALESCE PARTITION` can be used with a table that is partitioned by `HASH` or `KEY` to reduce the number of partitions by *`number`*. Suppose that you have created table `t2` as follows:

```
CREATE TABLE t2 (
    name VARCHAR (30),
    started DATE
)
PARTITION BY HASH( YEAR(started) )
```

```
PARTITIONS 6;
```

You can reduce the number of partitions used by `t2` from 6 to 4 using the following statement:

```
ALTER TABLE t2 COALESCE PARTITION 2;
```

The data contained in the last *`number`* partitions is merged into the remaining partitions. In this case, partitions 4 and 5 will be merged into the first 4 partitions (the partitions numbered 0, 1, 2, and 3).

To change some but not all the partitions used by a partitioned table, you can use `REORGANIZE PARTITION`. This statement can be used in several ways:

- To merge a set of partitions into a single partition. This can be done by naming several partitions in the *`partition_names`* list and supplying a single definition for *`partition_definition`*.
- To split an existing partition into several partitions. You can accomplish this by naming a single partition for *`partition_names`* and providing multiple *`partition_definitions`*.
- To change the ranges for a subset of partitions defined using `VALUES LESS THAN` or the value lists for a subset of partitions defined using `VALUES IN`.
- To move a partition from one tablespace to another. For an example, see Moving Table Partitions Between Tablespaces Using ALTER TABLE.

Note

For partitions that have not been explicitly named, MySQL automatically provides the default names `p0`, `p1`, `p2`, and so on. The same is true with regard to subpartitions.

For more detailed information about and examples of `ALTER TABLE ... REORGANIZE PARTITION` statements, see Section 22.3.1, “Management of RANGE and LIST Partitions”.

- It is possible to exchange a table partition or subpartition with a table using the `ALTER TABLE ... EXCHANGE PARTITION` statement—that is, to move any existing rows in the partition or subpartition to the nonpartitioned table, and any existing rows in the nonpartitioned table to the table partition or subpartition.

  For usage information and examples, see Section 22.3.3, “Exchanging Partitions and Subpartitions with Tables”.

- Several additional options provide partition maintenance and repair functionality analogous to that implemented for nonpartitioned tables by statements such as `CHECK TABLE` and `REPAIR TABLE` (which are also supported for partitioned tables; see Section 13.7.2, “Table Maintenance Statements” for more information). These include `ANALYZE PARTITION`, `CHECK PARTITION`, `OPTIMIZE PARTITION`, `REBUILD PARTITION`, and `REPAIR PARTITION`. Each of these options takes a *`partition_names`* clause consisting of one or more names of partitions, separated by commas. The partitions must already exist in the target table. You can also use the `ALL` keyword in place of *`partition_names`*, in which case the statement acts on all partitions in the table. For more information and examples, see Section 22.3.4, “Maintenance of Partitions”.

  `InnoDB` does not currently support per-partition optimization; `ALTER TABLE ... OPTIMIZE PARTITION` causes the entire table to rebuilt and analyzed, and an appropriate warning to be issued. (Bug #11751825, Bug #42822) To work around this problem, use `ALTER TABLE ... REBUILD PARTITION` and `ALTER TABLE ... ANALYZE PARTITION` instead.

  The `ANALYZE PARTITION`, `CHECK PARTITION`, `OPTIMIZE PARTITION`, and `REPAIR PARTITION` options are not supported for tables which are not partitioned.

- `REMOVE PARTITIONING` enables you to remove a table's partitioning without otherwise affecting the table or its data. This option can be combined with other `ALTER TABLE` options such as those used to add, drop, or rename columns or indexes.
- Using the `ENGINE` option with `ALTER TABLE` changes the storage engine used by the table without affecting the partitioning. The target storage engine must provide its own partitioning handler. Only the `InnoDB` and `NDB` storage engines have native partitioning handlers; `NDB` is not currently supported in MySQL 8.0.

It is possible for an `ALTER TABLE` statement to contain a `PARTITION BY` or `REMOVE PARTITIONING` clause in an addition to other alter specifications, but the `PARTITION BY` or `REMOVE PARTITIONING` clause must be specified last after any other specifications.

The `ADD PARTITION`, `DROP PARTITION`, `COALESCE PARTITION`, `REORGANIZE PARTITION`, `ANALYZE PARTITION`, `CHECK PARTITION`, and `REPAIR PARTITION` options cannot be combined with other alter specifications in a single `ALTER TABLE`, since the options just listed act on individual partitions. For more information, see Section 13.1.7.1, “ALTER TABLE Partition Operations”.

Only a single instance of any one of the following options can be used in a given `ALTER TABLE` statement: `PARTITION BY`, `ADD PARTITION`, `DROP PARTITION`, `TRUNCATE PARTITION`, `EXCHANGE PARTITION`, `REORGANIZE PARTITION`, or `COALESCE PARTITION`, `ANALYZE PARTITION`, `CHECK PARTITION`, `OPTIMIZE PARTITION`, `REBUILD PARTITION`, `REMOVE PARTITIONING`.

For example, the following two statements are invalid:

```
ALTER TABLE t1 ANALYZE PARTITION p1, ANALYZE PARTITION p2;

ALTER TABLE t1 ANALYZE PARTITION p1, CHECK PARTITION p2;
```

In the first case, you can analyze partitions `p1` and `p2` of table `t1` concurrently using a single statement with a single `ANALYZE PARTITION` option that lists both of the partitions to be analyzed, like this:

```
ALTER TABLE t1 ANALYZE PARTITION p1, p2;
```

In the second case, it is not possible to perform `ANALYZE` and `CHECK` operations on different partitions of the same table concurrently. Instead, you must issue two separate statements, like this:

```
ALTER TABLE t1 ANALYZE PARTITION p1;
ALTER TABLE t1 CHECK PARTITION p2;
```

`REBUILD` operations are currently unsupported for subpartitions. The `REBUILD` keyword is expressly disallowed with subpartitions, and causes `ALTER TABLE` to fail with an error if so used.

`CHECK PARTITION` and `REPAIR PARTITION` operations fail when the partition to be checked or repaired contains any duplicate key errors.

For more information about these statements, see Section 22.3.4, “Maintenance of Partitions”.

13.1.7.2 ALTER TABLE and Generated Columns

`ALTER TABLE` operations permitted for generated columns are `ADD`, `MODIFY`, and `CHANGE`.

- Generated columns can be added.
- The data type and expression of generated columns can be modified.
- Generated columns can be renamed or dropped, if no other column refers to them.
- Virtual generated columns cannot be altered to stored generated columns, or vice versa. To work around this, drop the column, then add it with the new definition.
- Nongenerated columns can be altered to stored but not virtual generated columns.

- Stored but not virtual generated columns can be altered to nongenerated columns. The stored generated values become the values of the nongenerated column.
- `ADD COLUMN` is not an in-place operation for stored columns (done without using a temporary table) because the expression must be evaluated by the server. For stored columns, indexing changes are done in place, and expression changes are not done in place. Changes to column comments are done in place.
- For non-partitioned tables, `ADD COLUMN` and `DROP COLUMN` are in-place operations for virtual columns. However, adding or dropping a virtual column cannot be performed in place in combination with other `ALTER TABLE` operations.

  For partitioned tables, `ADD COLUMN` and `DROP COLUMN` are not in-place operations for virtual columns.
- `InnoDB` supports secondary indexes on virtual generated columns. Adding or dropping a secondary index on a virtual generated column is an in-place operation. For more information, see Section 13.1.15.9, “Secondary Indexes and Generated Columns”.
- When a `VIRTUAL` generated column is added to a table or modified, it is not ensured that data being calculated by the generated column expression will not be out of range for the column. This can lead to inconsistent data being returned and unexpectedly failed statements. To permit control over whether validation occurs for such columns, `ALTER TABLE` supports `WITHOUT VALIDATION` and `WITH VALIDATION` clauses:
  - With `WITHOUT VALIDATION` (the default if neither clause is specified), an in-place operation is performed (if possible), data integrity is not checked, and the statement finishes more quickly. However, later reads from the table might report warnings or errors for the column if values are out of range.
  - With `WITH VALIDATION`, `ALTER TABLE` copies the table. If an out-of-range or any other error occurs, the statement fails. Because a table copy is performed, the statement takes longer.

  `WITHOUT VALIDATION` and `WITH VALIDATION` are permitted only with `ADD COLUMN`, `CHANGE COLUMN`, and `MODIFY COLUMN` operations. An `ER_WRONG_USAGE` error occurs otherwise.
- If expression evaluation causes truncation or provides incorrect input to a function, the `ALTER TABLE` statement terminates with an error and the DDL operation is rejected.
- An `ALTER TABLE` statement that changes the default value of a column *`col_name`* may also change the value of a generated column expression that refers to the column using *`col_name`*, which may change the value of a generated column expression that refers to the column using `DEFAULT(`*`col_name`*`)`. For this reason, `ALTER TABLE` operations that change the definition of a column now cause a table rebuild if any generated column expression uses `DEFAULT()`.

13.1.7.3 ALTER TABLE Examples

Begin with a table `t1` that is created as shown here:

```
CREATE TABLE t1 (a INTEGER,b CHAR(10));
```

To rename the table from `t1` to `t2`:

```
ALTER TABLE t1 RENAME t2;
```

To change column `a` from `INTEGER` to `TINYINT NOT NULL` (leaving the name the same), and to change column `b` from `CHAR(10)` to `CHAR(20)` as well as renaming it from `b` to `c`:

```
ALTER TABLE t2 MODIFY a TINYINT NOT NULL, CHANGE b c CHAR(20);
```

To add a new `TIMESTAMP` column named `d`:

```
ALTER TABLE t2 ADD d TIMESTAMP;
```

To add an index on column `d` and a `UNIQUE` index on column `a`:

```
ALTER TABLE t2 ADD INDEX (d), ADD UNIQUE (a);
```

To remove column `c`:

```
ALTER TABLE t2 DROP COLUMN c;
```

To add a new `AUTO_INCREMENT` integer column named `c`:

```
ALTER TABLE t2 ADD c INT UNSIGNED NOT NULL AUTO_INCREMENT,
  ADD PRIMARY KEY (c);
```

We indexed `c` (as a `PRIMARY KEY`) because `AUTO_INCREMENT` columns must be indexed, and we declare `c` as `NOT NULL` because primary key columns cannot be `NULL`.

When you add an `AUTO_INCREMENT` column, column values are filled in with sequence numbers automatically. For `MyISAM` tables, you can set the first sequence number by executing `SET INSERT_ID=value` before `ALTER TABLE` or by using the `AUTO_INCREMENT=value` table option. See Section 5.1.5, “Server System Variables”.

With `MyISAM` tables, if you do not change the `AUTO_INCREMENT` column, the sequence number is not affected. If you drop an `AUTO_INCREMENT` column and then add another `AUTO_INCREMENT` column, the numbers are resequenced beginning with 1.

When replication is used, adding an `AUTO_INCREMENT` column to a table might not produce the same ordering of the rows on the slave and the master. This occurs because the order in which the rows are numbered depends on the specific storage engine used for the table and the order in which the rows were inserted. If it is important to have the same order on the master and slave, the rows must be ordered before assigning an `AUTO_INCREMENT` number. Assuming that you want to add an `AUTO_INCREMENT` column to the table `t1`, the following statements produce a new table `t2` identical to `t1` but with an `AUTO_INCREMENT` column:

```
CREATE TABLE t2 (id INT AUTO_INCREMENT PRIMARY KEY)
SELECT * FROM t1 ORDER BY col1, col2;
```

This assumes that the table `t1` has columns `col1` and `col2`.

This set of statements will also produce a new table `t2` identical to `t1`, with the addition of an `AUTO_INCREMENT` column:

```
CREATE TABLE t2 LIKE t1;
ALTER TABLE t2 ADD id INT AUTO_INCREMENT PRIMARY KEY;
INSERT INTO t2 SELECT * FROM t1 ORDER BY col1, col2;
```

Important

To guarantee the same ordering on both master and slave, *all* columns of `t1` must be referenced in the `ORDER BY` clause.

Regardless of the method used to create and populate the copy having the `AUTO_INCREMENT` column, the final step is to drop the original table and then rename the copy:

```
DROP TABLE t1;
ALTER TABLE t2 RENAME t1;
```

13.1.8 ALTER VIEW Syntax

```
ALTER
    [ALGORITHM = {UNDEFINED | MERGE | TEMPTABLE}]
    [DEFINER = { user | CURRENT_USER }]
    [SQL SECURITY { DEFINER | INVOKER }]
    VIEW view_name [(column_list)]
    AS select_statement
    [WITH [CASCADED | LOCAL] CHECK OPTION]
```

This statement changes the definition of a view, which must exist. The syntax is similar to that for `CREATE VIEW` and the effect is the same as for `CREATE OR REPLACE VIEW`. See Section 13.1.18, “CREATE VIEW Syntax”. This statement requires the `CREATE VIEW` and `DROP` privileges for the view, and some privilege for each column referred to in the `SELECT` statement. `ALTER VIEW` is permitted only to the definer or users with the `SET_USER_ID` or `SUPER` privilege.

13.1.9 CREATE DATABASE Syntax

```
CREATE {DATABASE | SCHEMA} [IF NOT EXISTS] db_name
    [create_specification] ...

create_specification:
    [DEFAULT] CHARACTER SET [=] charset_name
  | [DEFAULT] COLLATE [=] collation_name
```

`CREATE DATABASE` creates a database with the given name. To use this statement, you need the `CREATE` privilege for the database. `CREATE SCHEMA` is a synonym for `CREATE DATABASE`.

An error occurs if the database exists and you did not specify `IF NOT EXISTS`.

In MySQL 8.0, `CREATE DATABASE` is not permitted within a session that has an active `LOCK TABLES` statement.

`create_specification` options specify database characteristics. Database characteristics are stored in the data dictionary. The `CHARACTER SET` clause specifies the default database character set. The `COLLATE` clause specifies the default database collation. Section 10.1, “Character Set Support”, discusses character set and collation names.

A database in MySQL is implemented as a directory containing files that correspond to tables in the database. Because there are no tables in a database when it is initially created, the `CREATE DATABASE` statement creates only a directory under the MySQL data directory. Rules for permissible database names are given in Section 9.2, “Schema Object Names”. If a database name contains special characters, the name for the database directory contains encoded versions of those characters as described in Section 9.2.3, “Mapping of Identifiers to File Names”.

Creating a database directory by manually creating a directory under the data directory (for example, with `mkdir`) is temporarily unsupported in MySQL 8.0.0.

You can also use the `mysqladmin` program to create databases. See Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”.

13.1.10 CREATE EVENT Syntax

```
CREATE
    [DEFINER = { user | CURRENT_USER }]
    EVENT
    [IF NOT EXISTS]
    event_name
    ON SCHEDULE schedule
    [ON COMPLETION [NOT] PRESERVE]
    [ENABLE | DISABLE | DISABLE ON SLAVE]
    [COMMENT 'comment']
    DO event_body;

schedule:
```

```
    AT timestamp [+ INTERVAL interval] ...
  | EVERY interval
    [STARTS timestamp [+ INTERVAL interval] ...]
    [ENDS timestamp [+ INTERVAL interval] ...]

interval:
    quantity {YEAR | QUARTER | MONTH | DAY | HOUR | MINUTE |
              WEEK | SECOND | YEAR_MONTH | DAY_HOUR | DAY_MINUTE |
              DAY_SECOND | HOUR_MINUTE | HOUR_SECOND | MINUTE_SECOND}
```

This statement creates and schedules a new event. The event will not run unless the Event Scheduler is enabled. For information about checking Event Scheduler status and enabling it if necessary, see Section 23.4.2, “Event Scheduler Configuration”.

`CREATE EVENT` requires the `EVENT` privilege for the schema in which the event is to be created. It might also require the `SET_USER_ID` or `SUPER` privilege, depending on the `DEFINER` value, as described later in this section.

The minimum requirements for a valid `CREATE EVENT` statement are as follows:

- The keywords `CREATE EVENT` plus an event name, which uniquely identifies the event in a database schema.
- An `ON SCHEDULE` clause, which determines when and how often the event executes.
- A `DO` clause, which contains the SQL statement to be executed by an event.

This is an example of a minimal `CREATE EVENT` statement:

```
CREATE EVENT myevent
    ON SCHEDULE AT CURRENT_TIMESTAMP + INTERVAL 1 HOUR
    DO
      UPDATE myschema.mytable SET mycol = mycol + 1;
```

The previous statement creates an event named `myevent`. This event executes once—one hour following its creation—by running an SQL statement that increments the value of the `myschema.mytable` table's `mycol` column by 1.

The *`event_name`* must be a valid MySQL identifier with a maximum length of 64 characters. Event names are not case sensitive, so you cannot have two events named `myevent` and `MyEvent` in the same schema. In general, the rules governing event names are the same as those for names of stored routines. See Section 9.2, “Schema Object Names”.

An event is associated with a schema. If no schema is indicated as part of *`event_name`*, the default (current) schema is assumed. To create an event in a specific schema, qualify the event name with a schema using *`schema_name.event_name`* syntax.

The `DEFINER` clause specifies the MySQL account to be used when checking access privileges at event execution time. If a *`user`* value is given, it should be a MySQL account specified as *`'user_name'@'host_name'`*, `CURRENT_USER`, or `CURRENT_USER()`. The default `DEFINER` value is the user who executes the `CREATE EVENT` statement. This is the same as specifying `DEFINER = CURRENT_USER` explicitly.

If you specify the `DEFINER` clause, these rules determine the valid `DEFINER` user values:

- If you do not have the `SET_USER_ID` or `SUPER` privilege, the only permitted *`user`* value is your own account, either specified literally or by using `CURRENT_USER`. You cannot set the definer to some other account.
- If you have the `SET_USER_ID` or `SUPER` privilege, you can specify any syntactically valid account name. If the account does not exist, a warning is generated.
- Although it is possible to create an event with a nonexistent `DEFINER` account, an error occurs at event execution time if the account does not exist.

For more information about event security, see Section 23.6, “Access Control for Stored Programs and Views”.

Within an event, the `CURRENT_USER()` function returns the account used to check privileges at event execution time, which is the `DEFINER` user. For information about user auditing within events, see Section 6.3.13, “SQL-Based MySQL Account Activity Auditing”.

`IF NOT EXISTS` has the same meaning for `CREATE EVENT` as for `CREATE TABLE`: If an event named *`event_name`* already exists in the same schema, no action is taken, and no error results. (However, a warning is generated in such cases.)

The `ON SCHEDULE` clause determines when, how often, and for how long the *`event_body`* defined for the event repeats. This clause takes one of two forms:

- `AT` *`timestamp`* is used for a one-time event. It specifies that the event executes one time only at the date and time given by *`timestamp`*, which must include both the date and time, or must be an expression that resolves to a datetime value. You may use a value of either the `DATETIME` or `TIMESTAMP` type for this purpose. If the date is in the past, a warning occurs, as shown here:

```
mysql> SELECT NOW();
+---------------------+
| NOW()               |
+---------------------+
| 2006-02-10 23:59:01 |
+---------------------+
1 row in set (0.04 sec)

mysql> CREATE EVENT e_totals
    ->     ON SCHEDULE AT '2006-02-10 23:59:00'
    ->     DO INSERT INTO test.totals VALUES (NOW());
Query OK, 0 rows affected, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1588
Message: Event execution time is in the past and ON COMPLETION NOT
         PRESERVE is set. The event was dropped immediately after
         creation.
```

  `CREATE EVENT` statements which are themselves invalid—for whatever reason—fail with an error.

  You may use `CURRENT_TIMESTAMP` to specify the current date and time. In such a case, the event acts as soon as it is created.

  To create an event which occurs at some point in the future relative to the current date and time—such as that expressed by the phrase “three weeks from now”—you can use the optional clause `+ INTERVAL` *`interval`*. The *`interval`* portion consists of two parts, a quantity and a unit of time, and follows the same syntax rules that govern intervals used in the `DATE_ADD()` function (see Section 12.7, “Date and Time Functions”. The units keywords are also the same, except that you cannot use any units involving microseconds when defining an event. With some interval types, complex time units may be used. For example, “two minutes and ten seconds” can be expressed as `+ INTERVAL '2:10' MINUTE_SECOND`.

  You can also combine intervals. For example, `AT CURRENT_TIMESTAMP + INTERVAL 3 WEEK + INTERVAL 2 DAY` is equivalent to “three weeks and two days from now”. Each portion of such a clause must begin with `+ INTERVAL`.

- To repeat actions at a regular interval, use an `EVERY` clause. The `EVERY` keyword is followed by an *`interval`* as described in the previous discussion of the `AT` keyword. (`+ INTERVAL` is *not* used with `EVERY`.) For example, `EVERY 6 WEEK` means “every six weeks”.

  Although `+ INTERVAL` clauses are not permitted in an `EVERY` clause, you can use the same complex time units permitted in a `+ INTERVAL`.

An EVERY clause may contain an optional STARTS clause. STARTS is followed by a *timestamp* value that indicates when the action should begin repeating, and may also use + INTERVAL *interval* to specify an amount of time "from now". For example, EVERY 3 MONTH STARTS CURRENT_TIMESTAMP + INTERVAL 1 WEEK means "every three months, beginning one week from now". Similarly, you can express "every two weeks, beginning six hours and fifteen minutes from now" as EVERY 2 WEEK STARTS CURRENT_TIMESTAMP + INTERVAL '6:15' HOUR_MINUTE. Not specifying STARTS is the same as using STARTS CURRENT_TIMESTAMP—that is, the action specified for the event begins repeating immediately upon creation of the event.

An EVERY clause may contain an optional ENDS clause. The ENDS keyword is followed by a *timestamp* value that tells MySQL when the event should stop repeating. You may also use + INTERVAL *interval* with ENDS; for instance, EVERY 12 HOUR STARTS CURRENT_TIMESTAMP + INTERVAL 30 MINUTE ENDS CURRENT_TIMESTAMP + INTERVAL 4 WEEK is equivalent to "every twelve hours, beginning thirty minutes from now, and ending four weeks from now". Not using ENDS means that the event continues executing indefinitely.

ENDS supports the same syntax for complex time units as STARTS does.

You may use STARTS, ENDS, both, or neither in an EVERY clause.

If a repeating event does not terminate within its scheduling interval, the result may be multiple instances of the event executing simultaneously. If this is undesirable, you should institute a mechanism to prevent simultaneous instances. For example, you could use the GET_LOCK() function, or row or table locking.

The ON SCHEDULE clause may use expressions involving built-in MySQL functions and user variables to obtain any of the *timestamp* or *interval* values which it contains. You may not use stored functions or user-defined functions in such expressions, nor may you use any table references; however, you may use SELECT FROM DUAL. This is true for both CREATE EVENT and ALTER EVENT statements. References to stored functions, user-defined functions, and tables in such cases are specifically not permitted, and fail with an error (see Bug #22830).

Times in the ON SCHEDULE clause are interpreted using the current session time_zone value. This becomes the event time zone; that is, the time zone that is used for event scheduling and is in effect within the event as it executes. These times are converted to UTC and stored along with the event time zone in the mysql.event table. This enables event execution to proceed as defined regardless of any subsequent changes to the server time zone or daylight saving time effects. For additional information about representation of event times, see Section 23.4.4, "Event Metadata". See also Section 13.7.5.18, "SHOW EVENTS Syntax", and Section 24.8, "The INFORMATION_SCHEMA EVENTS Table".

Normally, once an event has expired, it is immediately dropped. You can override this behavior by specifying ON COMPLETION PRESERVE. Using ON COMPLETION NOT PRESERVE merely makes the default nonpersistent behavior explicit.

You can create an event but prevent it from being active using the DISABLE keyword. Alternatively, you can use ENABLE to make explicit the default status, which is active. This is most useful in conjunction with ALTER EVENT (see Section 13.1.2, "ALTER EVENT Syntax").

A third value may also appear in place of ENABLE or DISABLE; DISABLE ON SLAVE is set for the status of an event on a replication slave to indicate that the event was created on the master and replicated to the slave, but is not executed on the slave. See Section 18.4.1.12, "Replication of Invoked Features".

You may supply a comment for an event using a COMMENT clause. *comment* may be any string of up to 64 characters that you wish to use for describing the event. The comment text, being a string literal, must be surrounded by quotation marks.

The DO clause specifies an action carried by the event, and consists of an SQL statement. Nearly any valid MySQL statement that can be used in a stored routine can also be used as the action

statement for a scheduled event. (See Section C.1, “Restrictions on Stored Programs”.) For example, the following event `e_hourly` deletes all rows from the `sessions` table once per hour, where this table is part of the `site_activity` schema:

```
CREATE EVENT e_hourly
    ON SCHEDULE
      EVERY 1 HOUR
    COMMENT 'Clears out sessions table each hour.'
    DO
      DELETE FROM site_activity.sessions;
```

MySQL stores the `sql_mode` system variable setting in effect when an event is created or altered, and always executes the event with this setting in force, *regardless of the current server SQL mode when the event begins executing*.

A `CREATE EVENT` statement that contains an `ALTER EVENT` statement in its `DO` clause appears to succeed; however, when the server attempts to execute the resulting scheduled event, the execution fails with an error.

Note

Statements such as `SELECT` or `SHOW` that merely return a result set have no effect when used in an event; the output from these is not sent to the MySQL Monitor, nor is it stored anywhere. However, you can use statements such as `SELECT ... INTO` and `INSERT INTO ... SELECT` that store a result. (See the next example in this section for an instance of the latter.)

The schema to which an event belongs is the default schema for table references in the `DO` clause. Any references to tables in other schemas must be qualified with the proper schema name.

As with stored routines, you can use compound-statement syntax in the `DO` clause by using the `BEGIN` and `END` keywords, as shown here:

```
delimiter |

CREATE EVENT e_daily
    ON SCHEDULE
      EVERY 1 DAY
    COMMENT 'Saves total number of sessions then clears the table each day'
    DO
      BEGIN
        INSERT INTO site_activity.totals (time, total)
          SELECT CURRENT_TIMESTAMP, COUNT(*)
            FROM site_activity.sessions;
        DELETE FROM site_activity.sessions;
      END |

delimiter ;
```

This example uses the `delimiter` command to change the statement delimiter. See Section 23.1, “Defining Stored Programs”.

More complex compound statements, such as those used in stored routines, are possible in an event. This example uses local variables, an error handler, and a flow control construct:

```
delimiter |

CREATE EVENT e
    ON SCHEDULE
      EVERY 5 SECOND
    DO
      BEGIN
        DECLARE v INTEGER;
        DECLARE CONTINUE HANDLER FOR SQLEXCEPTION BEGIN END;
```

```
        SET v = 0;

        WHILE v < 5 DO
          INSERT INTO t1 VALUES (0);
          UPDATE t2 SET s1 = s1 + 1;
          SET v = v + 1;
        END WHILE;
    END |

delimiter ;
```

There is no way to pass parameters directly to or from events; however, it is possible to invoke a stored routine with parameters within an event:

```
CREATE EVENT e_call_myproc
    ON SCHEDULE
      AT CURRENT_TIMESTAMP + INTERVAL 1 DAY
    DO CALL myproc(5, 27);
```

If an event's definer has the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege, the event can read and write global variables. As granting this privilege entails a potential for abuse, extreme care must be taken in doing so.

Generally, any statements that are valid in stored routines may be used for action statements executed by events. For more information about statements permissible within stored routines, see Section 23.2.1, “Stored Routine Syntax”. You can create an event as part of a stored routine, but an event cannot be created by another event.

13.1.11 CREATE FUNCTION Syntax

The `CREATE FUNCTION` statement is used to create stored functions and user-defined functions (UDFs):

- For information about creating stored functions, see Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”.
- For information about creating user-defined functions, see Section 13.7.3.1, “CREATE FUNCTION Syntax for User-Defined Functions”.

13.1.12 CREATE INDEX Syntax

```
CREATE [UNIQUE|FULLTEXT|SPATIAL] INDEX index_name
    [index_type]
    ON tbl_name (index_col_name,...)
    [index_option]
    [algorithm_option | lock_option] ...

index_col_name:
    col_name [(length)] [ASC | DESC]

index_option:
    KEY_BLOCK_SIZE [=] value
  | index_type
  | WITH PARSER parser_name
  | COMMENT 'string'
  | {VISIBLE | INVISIBLE}

index_type:
    USING {BTREE | HASH}

algorithm_option:
    ALGORITHM [=] {DEFAULT|INPLACE|COPY}

lock_option:
    LOCK [=] {DEFAULT|NONE|SHARED|EXCLUSIVE}
```

`CREATE INDEX` is mapped to an `ALTER TABLE` statement to create indexes. See Section 13.1.7, "ALTER TABLE Syntax". `CREATE INDEX` cannot be used to create a `PRIMARY KEY`; use `ALTER TABLE` instead. For more information about indexes, see Section 8.3.1, "How MySQL Uses Indexes".

Normally, you create all indexes on a table at the time the table itself is created with `CREATE TABLE`. See Section 13.1.15, "CREATE TABLE Syntax". This guideline is especially important for `InnoDB` tables, where the primary key determines the physical layout of rows in the data file. `CREATE INDEX` enables you to add indexes to existing tables.

A column list of the form `(col1, col2, ...)` creates a multiple-column index. Index key values are formed by concatenating the values of the given columns.

For string columns, indexes can be created that use only the leading part of column values, using *`col_name(length)`* syntax to specify an index prefix length:

- Prefixes can be specified for `CHAR`, `VARCHAR`, `BINARY`, and `VARBINARY` column indexes.
- Prefixes *must* be specified for `BLOB` and `TEXT` column indexes.
- Prefix limits are measured in bytes, whereas the prefix length in `CREATE TABLE`, `ALTER TABLE`, and `CREATE INDEX` statements is interpreted as number of characters for nonbinary string types (`CHAR`, `VARCHAR`, `TEXT`) and number of bytes for binary string types (`BINARY`, `VARBINARY`, `BLOB`). Take this into account when specifying a prefix length for a nonbinary string column that uses a multibyte character set.
- For spatial columns, prefix values cannot be given, as described later in this section.

The statement shown here creates an index using the first 10 characters of the `name` column (assuming that `name` has a nonbinary string type):

```
CREATE INDEX part_of_name ON customer (name(10));
```

If names in the column usually differ in the first 10 characters, this index should not be much slower than an index created from the entire `name` column. Also, using column prefixes for indexes can make the index file much smaller, which could save a lot of disk space and might also speed up `INSERT` operations.

Prefix support and lengths of prefixes (where supported) are storage engine dependent. For example, a prefix can be up to 767 bytes long for `InnoDB` tables that use the `REDUNDANT` or `COMPACT` row format. The length limit is raised to 3072 bytes for `InnoDB` tables that use the `DYNAMIC` or `COMPRESSED` row format. For `MyISAM` tables, the prefix limit is 1000 bytes.

A `UNIQUE` index creates a constraint such that all values in the index must be distinct. An error occurs if you try to add a new row with a key value that matches an existing row. For all engines, a `UNIQUE` index permits multiple `NULL` values for columns that can contain `NULL`. If you specify a prefix value for a column in a `UNIQUE` index, the column values must be unique within the prefix.

`FULLTEXT` indexes are supported only for `InnoDB` and `MyISAM` tables and can include only `CHAR`, `VARCHAR`, and `TEXT` columns. Indexing always happens over the entire column; column prefix indexing is not supported and any prefix length is ignored if specified. See Section 12.9, "Full-Text Search Functions", for details of operation.

The `MyISAM`, `InnoDB`, `NDB`, and `ARCHIVE` storage engines support spatial columns such as (`POINT` and `GEOMETRY`. (Section 11.5, "Extensions for Spatial Data", describes the spatial data types.) However, support for spatial column indexing varies among engines. Spatial and nonspatial indexes are available according to the following rules.

Spatial indexes (created using `SPATIAL INDEX`) have these characteristics:

- Available only for `InnoDB` and `MyISAM` tables. Specifying `SPATIAL INDEX` for other storage engines results in an error.
- Indexed columns must be `NOT NULL`.

- Column prefix lengths are prohibited. The full width of each column is indexed.

Characteristics of nonspatial indexes (created with `INDEX`, `UNIQUE`, or `PRIMARY KEY`):

- Permitted for any storage engine that supports spatial columns except `ARCHIVE`.
- Columns can be `NULL` unless the index is a primary key.
- For each spatial column in a non-`SPATIAL` index except `POINT` columns, a column prefix length must be specified. (This is the same requirement as for indexed `BLOB` columns.) The prefix length is given in bytes.
- The index type for a non-`SPATIAL` index depends on the storage engine. Currently, B-tree is used.
- You can add an index on a column that can have `NULL` values only for `InnoDB`, `MyISAM`, and `MEMORY` tables.
- You can add an index on a `BLOB` or `TEXT` column only for using the `InnoDB` and `MyISAM` tables.
- When the `innodb_stats_persistent` setting is enabled, run the `ANALYZE TABLE` statement for an `InnoDB` table after creating an index on that table.

`InnoDB` supports secondary indexes on virtual columns. For more information, see Section 13.1.15.9, “Secondary Indexes and Generated Columns”.

An *`index_col_name`* specification can end with `ASC` or `DESC` to specify whether index values are stored in ascending or descending order. The default is ascending if no order specifier is given. `ASC` and `DESC` are not permitted for `HASH` indexes.

Following the index column list, index options can be given. An *`index_option`* value can be any of the following:

- `KEY_BLOCK_SIZE [=]` *`value`*

  For `MyISAM` tables, `KEY_BLOCK_SIZE` optionally specifies the size in bytes to use for index key blocks. The value is treated as a hint; a different size could be used if necessary. A `KEY_BLOCK_SIZE` value specified for an individual index definition overrides a table-level `KEY_BLOCK_SIZE` value.

  `KEY_BLOCK_SIZE` is not supported at the index level for `InnoDB` tables. See Section 13.1.15, “CREATE TABLE Syntax”.

- *`index_type`*

  Some storage engines permit you to specify an index type when creating an index. For example:

```
CREATE TABLE lookup (id INT) ENGINE = MEMORY;
CREATE INDEX id_index ON lookup (id) USING BTREE;
```

  Table 13.1, “Index Types Per Storage Engine” shows the permissible index type values supported by different storage engines. Where multiple index types are listed, the first one is the default when no index type specifier is given. Storage engines not listed in the table do not support an *`index_type`* clause in index definitions.

Table 13.1 Index Types Per Storage Engine

Storage Engine	Permissible Index Types
`InnoDB`	`BTREE`
`MyISAM`	`BTREE`
`MEMORY`/`HEAP`	`HASH`, `BTREE`
`NDB`	`HASH`, `BTREE` (see note in text)

The *index_type* clause cannot be used for FULLTEXT INDEX or SPATIAL INDEX specifications. Full-text index implementation is storage engine dependent. Spatial indexes are implemented as R-tree indexes.

If you specify an index type that is not valid for a given storage engine, but another index type is available that the engine can use without affecting query results, the engine uses the available type. The parser recognizes RTREE as a type name, but currently this cannot be specified for any storage engine.

Note

Use of the *index_type* option before the ON *tbl_name* clause is deprecated; support for use of the option in this position will be removed in a future MySQL release. If an *index_type* option is given in both the earlier and later positions, the final option applies.

TYPE *type_name* is recognized as a synonym for USING *type_name*. However, USING is the preferred form.

For the storage engines that support an *index_type* option, Table 13.2, "Storage Engine Index Characteristics" shows some characteristics of index use.

Table 13.2 Storage Engine Index Characteristics

Storage Engine	Index Type	Index Class	Stores NULL Values	Permits Multiple NULL Values	IS NULL Scan Type	IS NOT NULL Scan Type
InnoDB	BTREE	Primary key	No	No	N/A	N/A
		Unique	Yes	Yes	Index	Index
		Key	Yes	Yes	Index	Index
	Inapplicable	FULLTEXT	Yes	Yes	Table	Table
	Inapplicable	SPATIAL	No	No	N/A	N/A
MyISAM	BTREE	Primary key	No	No	N/A	N/A
		Unique	Yes	Yes	Index	Index
		Key	Yes	Yes	Index	Index
	Inapplicable	FULLTEXT	Yes	Yes	Table	Table
	Inapplicable	SPATIAL	No	No	N/A	N/A
MEMORY	HASH	Primary key	No	No	N/A	N/A
		Unique	Yes	Yes	Index	Index
		Key	Yes	Yes	Index	Index
	BTREE	Primary	No	No	N/A	N/A
		Unique	Yes	Yes	Index	Index
		Key	Yes	Yes	Index	Index

- WITH PARSER *parser_name*

  This option can be used only with FULLTEXT indexes. It associates a parser plugin with the index if full-text indexing and searching operations need special handling. InnoDB and MyISAM support

full-text parser plugins. See Full-Text Parser Plugins and Section 28.2.4.4, “Writing Full-Text Parser Plugins” for more information.

- `COMMENT 'string'`

  Index definitions can include an optional comment of up to 1024 characters.

  The `MERGE_THRESHOLD` for index pages can be configured for individual indexes using the *`index_option`* `COMMENT` clause of the `CREATE INDEX` statement. For example:

```
CREATE TABLE t1 (id INT);
CREATE INDEX id_index ON t1 (id) COMMENT 'MERGE_THRESHOLD=40';
```

  If the page-full percentage for an index page falls below the `MERGE_THRESHOLD` value when a row is deleted or when a row is shortened by an update operation, `InnoDB` attempts to merge the index page with a neighboring index page. The default `MERGE_THRESHOLD` value is 50, which is the previously hardcoded value.

  `MERGE_THRESHOLD` can also be defined at the index level and table level using `CREATE TABLE` and `ALTER TABLE` statements. For more information, see Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

- `VISIBLE`, `INVISIBLE`

  Modify index visibility. Indexes are visible by default. An invisible index is not used by the optimizer. Modification of index visibility applies to indexes other than primary keys (either explicit or implicit). This feature is storage engine neutral (supported for any engine). For more information, see Section 8.3.11, “Invisible Indexes”.

`ALGORITHM` and `LOCK` clauses may be given to influence the table copying method and level of concurrency for reading and writing the table while its indexes are being modified. They have the same meaning as for the `ALTER TABLE` statement. For more information, see Section 13.1.7, “ALTER TABLE Syntax”

13.1.13 CREATE PROCEDURE and CREATE FUNCTION Syntax

```
CREATE
    [DEFINER = { user | CURRENT_USER }]
    PROCEDURE sp_name ([proc_parameter[,...]])
    [characteristic ...] routine_body

CREATE
    [DEFINER = { user | CURRENT_USER }]
    FUNCTION sp_name ([func_parameter[,...]])
    RETURNS type
    [characteristic ...] routine_body

proc_parameter:
    [ IN | OUT | INOUT ] param_name type

func_parameter:
    param_name type

type:
    Any valid MySQL data type

characteristic:
    COMMENT 'string'
  | LANGUAGE SQL
  | [NOT] DETERMINISTIC
  | { CONTAINS SQL | NO SQL | READS SQL DATA | MODIFIES SQL DATA }
  | SQL SECURITY { DEFINER | INVOKER }

routine_body:
```

```
    Valid SQL routine statement
```

These statements create stored routines. By default, a routine is associated with the default database. To associate the routine explicitly with a given database, specify the name as *db_name.sp_name* when you create it.

The CREATE FUNCTION statement is also used in MySQL to support UDFs (user-defined functions). See Section 28.4, “Adding New Functions to MySQL”. A UDF can be regarded as an external stored function. Stored functions share their namespace with UDFs. See Section 9.2.4, “Function Name Parsing and Resolution”, for the rules describing how the server interprets references to different kinds of functions.

To invoke a stored procedure, use the CALL statement (see Section 13.2.1, “CALL Syntax”). To invoke a stored function, refer to it in an expression. The function returns a value during expression evaluation.

CREATE PROCEDURE and CREATE FUNCTION require the CREATE ROUTINE privilege. They might also require the SET_USER_ID or SUPER privilege, depending on the DEFINER value, as described later in this section. If binary logging is enabled, CREATE FUNCTION might require the SUPER privilege, as described in Section 23.7, “Binary Logging of Stored Programs”.

By default, MySQL automatically grants the ALTER ROUTINE and EXECUTE privileges to the routine creator. This behavior can be changed by disabling the automatic_sp_privileges system variable. See Section 23.2.2, “Stored Routines and MySQL Privileges”.

The DEFINER and SQL SECURITY clauses specify the security context to be used when checking access privileges at routine execution time, as described later in this section.

If the routine name is the same as the name of a built-in SQL function, a syntax error occurs unless you use a space between the name and the following parenthesis when defining the routine or invoking it later. For this reason, avoid using the names of existing SQL functions for your own stored routines.

The IGNORE_SPACE SQL mode applies to built-in functions, not to stored routines. It is always permissible to have spaces after a stored routine name, regardless of whether IGNORE_SPACE is enabled.

The parameter list enclosed within parentheses must always be present. If there are no parameters, an empty parameter list of () should be used. Parameter names are not case sensitive.

Each parameter is an IN parameter by default. To specify otherwise for a parameter, use the keyword OUT or INOUT before the parameter name.

Note

Specifying a parameter as IN, OUT, or INOUT is valid only for a PROCEDURE. For a FUNCTION, parameters are always regarded as IN parameters.

An IN parameter passes a value into a procedure. The procedure might modify the value, but the modification is not visible to the caller when the procedure returns. An OUT parameter passes a value from the procedure back to the caller. Its initial value is NULL within the procedure, and its value is visible to the caller when the procedure returns. An INOUT parameter is initialized by the caller, can be modified by the procedure, and any change made by the procedure is visible to the caller when the procedure returns.

For each OUT or INOUT parameter, pass a user-defined variable in the CALL statement that invokes the procedure so that you can obtain its value when the procedure returns. If you are calling the procedure from within another stored procedure or function, you can also pass a routine parameter or local routine variable as an IN or INOUT parameter.

Routine parameters cannot be referenced in statements prepared within the routine; see Section C.1, “Restrictions on Stored Programs”.

The following example shows a simple stored procedure that uses an OUT parameter:

```
mysql> delimiter //

mysql> CREATE PROCEDURE simpleproc (OUT param1 INT)
    -> BEGIN
    ->   SELECT COUNT(*) INTO param1 FROM t;
    -> END//
Query OK, 0 rows affected (0.00 sec)

mysql> delimiter ;

mysql> CALL simpleproc(@a);
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @a;
+------+
| @a   |
+------+
| 3    |
+------+
1 row in set (0.00 sec)
```

The example uses the `mysql` client `delimiter` command to change the statement delimiter from `;` to `//` while the procedure is being defined. This enables the `;` delimiter used in the procedure body to be passed through to the server rather than being interpreted by `mysql` itself. See Section 23.1, “Defining Stored Programs”.

The `RETURNS` clause may be specified only for a `FUNCTION`, for which it is mandatory. It indicates the return type of the function, and the function body must contain a `RETURN` *`value`* statement. If the `RETURN` statement returns a value of a different type, the value is coerced to the proper type. For example, if a function specifies an `ENUM` or `SET` value in the `RETURNS` clause, but the `RETURN` statement returns an integer, the value returned from the function is the string for the corresponding `ENUM` member of set of `SET` members.

The following example function takes a parameter, performs an operation using an SQL function, and returns the result. In this case, it is unnecessary to use `delimiter` because the function definition contains no internal `;` statement delimiters:

```
mysql> CREATE FUNCTION hello (s CHAR(20))
mysql> RETURNS CHAR(50) DETERMINISTIC
    -> RETURN CONCAT('Hello, ',s,'!');
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT hello('world');
+----------------+
| hello('world') |
+----------------+
| Hello, world!  |
+----------------+
1 row in set (0.00 sec)
```

Parameter types and function return types can be declared to use any valid data type. The `COLLATE` attribute can be used if preceded by the `CHARACTER SET` attribute.

The *`routine_body`* consists of a valid SQL routine statement. This can be a simple statement such as `SELECT` or `INSERT`, or a compound statement written using `BEGIN` and `END`. Compound statements can contain declarations, loops, and other control structure statements. The syntax for these statements is described in Section 13.6, “Compound-Statement Syntax”.

MySQL permits routines to contain DDL statements, such as `CREATE` and `DROP`. MySQL also permits stored procedures (but not stored functions) to contain SQL transaction statements such as `COMMIT`. Stored functions may not contain statements that perform explicit or implicit commit or rollback. Support for these statements is not required by the SQL standard, which states that each DBMS vendor may decide whether to permit them.

Statements that return a result set can be used within a stored procedure but not within a stored function. This prohibition includes `SELECT` statements that do not have an `INTO` *`var_list`* clause

and other statements such as `SHOW`, `EXPLAIN`, and `CHECK TABLE`. For statements that can be determined at function definition time to return a result set, a `Not allowed to return a result set from a function` error occurs (`ER_SP_NO_RETSET`). For statements that can be determined only at runtime to return a result set, a `PROCEDURE %s can't return a result set in the given context` error occurs (`ER_SP_BADSELECT`).

`USE` statements within stored routines are not permitted. When a routine is invoked, an implicit `USE` *`db_name`* is performed (and undone when the routine terminates). The causes the routine to have the given default database while it executes. References to objects in databases other than the routine default database should be qualified with the appropriate database name.

For additional information about statements that are not permitted in stored routines, see Section C.1, “Restrictions on Stored Programs”.

For information about invoking stored procedures from within programs written in a language that has a MySQL interface, see Section 13.2.1, “CALL Syntax”.

MySQL stores the `sql_mode` system variable setting in effect when a routine is created or altered, and always executes the routine with this setting in force, *regardless of the current server SQL mode when the routine begins executing*.

The switch from the SQL mode of the invoker to that of the routine occurs after evaluation of arguments and assignment of the resulting values to routine parameters. If you define a routine in strict SQL mode but invoke it in nonstrict mode, assignment of arguments to routine parameters does not take place in strict mode. If you require that expressions passed to a routine be assigned in strict SQL mode, you should invoke the routine with strict mode in effect.

The `COMMENT` characteristic is a MySQL extension, and may be used to describe the stored routine. This information is displayed by the `SHOW CREATE PROCEDURE` and `SHOW CREATE FUNCTION` statements.

The `LANGUAGE` characteristic indicates the language in which the routine is written. The server ignores this characteristic; only SQL routines are supported.

A routine is considered “deterministic” if it always produces the same result for the same input parameters, and “not deterministic” otherwise. If neither `DETERMINISTIC` nor `NOT DETERMINISTIC` is given in the routine definition, the default is `NOT DETERMINISTIC`. To declare that a function is deterministic, you must specify `DETERMINISTIC` explicitly.

Assessment of the nature of a routine is based on the “honesty” of the creator: MySQL does not check that a routine declared `DETERMINISTIC` is free of statements that produce nondeterministic results. However, misdeclaring a routine might affect results or affect performance. Declaring a nondeterministic routine as `DETERMINISTIC` might lead to unexpected results by causing the optimizer to make incorrect execution plan choices. Declaring a deterministic routine as `NONDETERMINISTIC` might diminish performance by causing available optimizations not to be used.

If binary logging is enabled, the `DETERMINISTIC` characteristic affects which routine definitions MySQL accepts. See Section 23.7, “Binary Logging of Stored Programs”.

A routine that contains the `NOW()` function (or its synonyms) or `RAND()` is nondeterministic, but it might still be replication-safe. For `NOW()`, the binary log includes the timestamp and replicates correctly. `RAND()` also replicates correctly as long as it is called only a single time during the execution of a routine. (You can consider the routine execution timestamp and random number seed as implicit inputs that are identical on the master and slave.)

Several characteristics provide information about the nature of data use by the routine. In MySQL, these characteristics are advisory only. The server does not use them to constrain what kinds of statements a routine will be permitted to execute.

- `CONTAINS SQL` indicates that the routine does not contain statements that read or write data. This is the default if none of these characteristics is given explicitly. Examples of such statements are `SET @x = 1` or `DO RELEASE_LOCK('abc')`, which execute but neither read nor write data.

- `NO SQL` indicates that the routine contains no SQL statements.
- `READS SQL DATA` indicates that the routine contains statements that read data (for example, `SELECT`), but not statements that write data.
- `MODIFIES SQL DATA` indicates that the routine contains statements that may write data (for example, `INSERT` or `DELETE`).

The `SQL SECURITY` characteristic can be `DEFINER` or `INVOKER` to specify the security context; that is, whether the routine executes using the privileges of the account named in the routine `DEFINER` clause or the user who invokes it. This account must have permission to access the database with which the routine is associated. The default value is `DEFINER`. The user who invokes the routine must have the `EXECUTE` privilege for it, as must the `DEFINER` account if the routine executes in definer security context.

The `DEFINER` clause specifies the MySQL account to be used when checking access privileges at routine execution time for routines that have the `SQL SECURITY DEFINER` characteristic.

If a *`user`* value is given for the `DEFINER` clause, it should be a MySQL account specified as *`'user_name'@'host_name'`*, `CURRENT_USER`, or `CURRENT_USER()`. The default `DEFINER` value is the user who executes the `CREATE PROCEDURE` or `CREATE FUNCTION` statement. This is the same as specifying `DEFINER = CURRENT_USER` explicitly.

If you specify the `DEFINER` clause, these rules determine the valid `DEFINER` user values:

- If you do not have the `SET_USER_ID` or `SUPER` privilege, the only permitted *`user`* value is your own account, either specified literally or by using `CURRENT_USER`. You cannot set the definer to some other account.
- If you have the `SET_USER_ID` or `SUPER` privilege, you can specify any syntactically valid account name. If the account does not exist, a warning is generated.
- Although it is possible to create a routine with a nonexistent `DEFINER` account, an error occurs at routine execution time if the `SQL SECURITY` value is `DEFINER` but the definer account does not exist.

For more information about stored routine security, see Section 23.6, “Access Control for Stored Programs and Views”.

Within a stored routine that is defined with the `SQL SECURITY DEFINER` characteristic, `CURRENT_USER` returns the routine's `DEFINER` value. For information about user auditing within stored routines, see Section 6.3.13, “SQL-Based MySQL Account Activity Auditing”.

Consider the following procedure, which displays a count of the number of MySQL accounts listed in the `mysql.user` table:

```
CREATE DEFINER = 'admin'@'localhost' PROCEDURE account_count()
BEGIN
  SELECT 'Number of accounts:', COUNT(*) FROM mysql.user;
END;
```

The procedure is assigned a `DEFINER` account of `'admin'@'localhost'` no matter which user defines it. It executes with the privileges of that account no matter which user invokes it (because the default security characteristic is `DEFINER`). The procedure succeeds or fails depending on whether invoker has the `EXECUTE` privilege for it and `'admin'@'localhost'` has the `SELECT` privilege for the `mysql.user` table.

Now suppose that the procedure is defined with the `SQL SECURITY INVOKER` characteristic:

```
CREATE DEFINER = 'admin'@'localhost' PROCEDURE account_count()
SQL SECURITY INVOKER
BEGIN
```

```
  SELECT 'Number of accounts:', COUNT(*) FROM mysql.user;
END;
```

The procedure still has a `DEFINER` of `'admin'@'localhost'`, but in this case, it executes with the privileges of the invoking user. Thus, the procedure succeeds or fails depending on whether the invoker has the `EXECUTE` privilege for it and the `SELECT` privilege for the `mysql.user` table.

The server handles the data type of a routine parameter, local routine variable created with `DECLARE`, or function return value as follows:

- Assignments are checked for data type mismatches and overflow. Conversion and overflow problems result in warnings, or errors in strict SQL mode.
- Only scalar values can be assigned. For example, a statement such as `SET x = (SELECT 1, 2)` is invalid.
- For character data types, if there is a `CHARACTER SET` attribute in the declaration, the specified character set and its default collation is used. If the `COLLATE` attribute is also present, that collation is used rather than the default collation.

  If `CHARACTER SET` and `COLLATE` attributes are not present, the database character set and collation in effect at routine creation time are used. To avoid having the server use the database character set and collation, provide explicit `CHARACTER SET` and `COLLATE` attributes for character data parameters.

  If you change the database default character set or collation, stored routines that use the database defaults must be dropped and recreated so that they use the new defaults.

  The database character set and collation are given by the value of the `character_set_database` and `collation_database` system variables. For more information, see Section 10.1.3.3, “Database Character Set and Collation”.

13.1.14 CREATE SERVER Syntax

```
CREATE SERVER server_name
    FOREIGN DATA WRAPPER wrapper_name
    OPTIONS (option [, option] ...)

option:
  { HOST character-literal
  | DATABASE character-literal
  | USER character-literal
  | PASSWORD character-literal
  | SOCKET character-literal
  | OWNER character-literal
  | PORT numeric-literal }
```

This statement creates the definition of a server for use with the `FEDERATED` storage engine. The `CREATE SERVER` statement creates a new row in the `servers` table in the `mysql` database. This statement requires the `SUPER` privilege.

The *`server_name`* should be a unique reference to the server. Server definitions are global within the scope of the server, it is not possible to qualify the server definition to a specific database. *`server_name`* has a maximum length of 64 characters (names longer than 64 characters are silently truncated), and is case insensitive. You may specify the name as a quoted string.

The *`wrapper_name`* should be `mysql`, and may be quoted with single quotation marks. Other values for *`wrapper_name`* are not currently supported.

For each *`option`* you must specify either a character literal or numeric literal. Character literals are UTF-8, support a maximum length of 64 characters and default to a blank (empty) string. String literals are silently truncated to 64 characters. Numeric literals must be a number between 0 and 9999, default value is 0.

Note

The `OWNER` option is currently not applied, and has no effect on the ownership or operation of the server connection that is created.

The `CREATE SERVER` statement creates an entry in the `mysql.servers` table that can later be used with the `CREATE TABLE` statement when creating a `FEDERATED` table. The options that you specify will be used to populate the columns in the `mysql.servers` table. The table columns are `Server_name`, `Host`, `Db`, `Username`, `Password`, `Port` and `Socket`.

For example:

```
CREATE SERVER s
FOREIGN DATA WRAPPER mysql
OPTIONS (USER 'Remote', HOST '192.168.1.106', DATABASE 'test');
```

Be sure to specify all options necessary to establish a connection to the server. The user name, host name, and database name are mandatory. Other options might be required as well, such as password.

The data stored in the table can be used when creating a connection to a `FEDERATED` table:

```
CREATE TABLE t (s1 INT) ENGINE=FEDERATED CONNECTION='s';
```

For more information, see Section 16.8, “The FEDERATED Storage Engine”.

`CREATE SERVER` causes an automatic commit.

In MySQL 8.0, `CREATE SERVER` is not written to the binary log, regardless of the logging format that is in use.

13.1.15 CREATE TABLE Syntax

```
CREATE [TEMPORARY] TABLE [IF NOT EXISTS] tbl_name
    (create_definition,...)
    [table_options]
    [partition_options]

CREATE [TEMPORARY] TABLE [IF NOT EXISTS] tbl_name
    [(create_definition,...)]
    [table_options]
    [partition_options]
    [IGNORE | REPLACE]
    [AS] query_expression

CREATE [TEMPORARY] TABLE [IF NOT EXISTS] tbl_name
    { LIKE old_tbl_name | (LIKE old_tbl_name) }

create_definition:
    col_name column_definition
  | [CONSTRAINT [symbol]] PRIMARY KEY [index_type] (index_col_name,...)
      [index_option] ...
  | {INDEX|KEY} [index_name] [index_type] (index_col_name,...)
      [index_option] ...
  | [CONSTRAINT [symbol]] UNIQUE [INDEX|KEY]
      [index_name] [index_type] (index_col_name,...)
      [index_option] ...
  | {FULLTEXT|SPATIAL} [INDEX|KEY] [index_name] (index_col_name,...)
      [index_option] ...
  | [CONSTRAINT [symbol]] FOREIGN KEY
      [index_name] (index_col_name,...) reference_definition
  | CHECK (expr)

column_definition:
    data_type [NOT NULL | NULL] [DEFAULT default_value]
      [AUTO_INCREMENT] [UNIQUE [KEY] | [PRIMARY] KEY]
      [COMMENT 'string']
      [COLUMN_FORMAT {FIXED|DYNAMIC|DEFAULT}]
```

```
      [reference_definition]
  | data_type [GENERATED ALWAYS] AS (expression)
      [VIRTUAL | STORED] [UNIQUE [KEY]] [COMMENT comment]
      [NOT NULL | NULL] [[PRIMARY] KEY]

data_type:
    BIT[(length)]
  | TINYINT[(length)] [UNSIGNED] [ZEROFILL]
  | SMALLINT[(length)] [UNSIGNED] [ZEROFILL]
  | MEDIUMINT[(length)] [UNSIGNED] [ZEROFILL]
  | INT[(length)] [UNSIGNED] [ZEROFILL]
  | INTEGER[(length)] [UNSIGNED] [ZEROFILL]
  | BIGINT[(length)] [UNSIGNED] [ZEROFILL]
  | REAL[(length,decimals)] [UNSIGNED] [ZEROFILL]
  | DOUBLE[(length,decimals)] [UNSIGNED] [ZEROFILL]
  | FLOAT[(length,decimals)] [UNSIGNED] [ZEROFILL]
  | DECIMAL[(length[,decimals])] [UNSIGNED] [ZEROFILL]
  | NUMERIC[(length[,decimals])] [UNSIGNED] [ZEROFILL]
  | DATE
  | TIME[(fsp)]
  | TIMESTAMP[(fsp)]
  | DATETIME[(fsp)]
  | YEAR
  | CHAR[(length)] [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | VARCHAR(length) [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | BINARY[(length)]
  | VARBINARY(length)
  | TINYBLOB
  | BLOB
  | MEDIUMBLOB
  | LONGBLOB
  | TINYTEXT [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | TEXT [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | MEDIUMTEXT [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | LONGTEXT [BINARY]
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | ENUM(value1,value2,value3,...)
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | SET(value1,value2,value3,...)
      [CHARACTER SET charset_name] [COLLATE collation_name]
  | JSON
  | spatial_type

index_col_name:
    col_name [(length)] [ASC | DESC]

index_type:
    USING {BTREE | HASH}

index_option:
    KEY_BLOCK_SIZE [=] value
  | index_type
  | WITH PARSER parser_name
  | COMMENT 'string'
  | {VISIBLE | INVISIBLE}

reference_definition:
    REFERENCES tbl_name (index_col_name,...)
      [MATCH FULL | MATCH PARTIAL | MATCH SIMPLE]
      [ON DELETE reference_option]
      [ON UPDATE reference_option]

reference_option:
    RESTRICT | CASCADE | SET NULL | NO ACTION | SET DEFAULT

table_options:
    table_option [[,] table_option] ...
```

```
table_option:
    ENGINE [=] engine_name
  | AUTO_INCREMENT [=] value
  | AVG_ROW_LENGTH [=] value
  | [DEFAULT] CHARACTER SET [=] charset_name
  | CHECKSUM [=] {0 | 1}
  | [DEFAULT] COLLATE [=] collation_name
  | COMMENT [=] 'string'
  | COMPRESSION [=] {'ZLIB'|'LZ4'|'NONE'}
  | CONNECTION [=] 'connect_string'
  | DATA DIRECTORY [=] 'absolute path to directory'
  | DELAY_KEY_WRITE [=] {0 | 1}
  | ENCRYPTION [=] {'Y' | 'N'}
  | INDEX DIRECTORY [=] 'absolute path to directory'
  | INSERT_METHOD [=] { NO | FIRST | LAST }
  | KEY_BLOCK_SIZE [=] value
  | MAX_ROWS [=] value
  | MIN_ROWS [=] value
  | PACK_KEYS [=] {0 | 1 | DEFAULT}
  | PASSWORD [=] 'string'
  | ROW_FORMAT [=] {DEFAULT|DYNAMIC|FIXED|COMPRESSED|REDUNDANT|COMPACT}
  | STATS_AUTO_RECALC [=] {DEFAULT|0|1}
  | STATS_PERSISTENT [=] {DEFAULT|0|1}
  | STATS_SAMPLE_PAGES [=] value
  | TABLESPACE tablespace_name
  | UNION [=] (tbl_name[,tbl_name]...)

partition_options:
    PARTITION BY
        { [LINEAR] HASH(expr)
        | [LINEAR] KEY [ALGORITHM={1|2}] (column_list)
        | RANGE{(expr) | COLUMNS(column_list)}
        | LIST{(expr) | COLUMNS(column_list)} }
    [PARTITIONS num]
    [SUBPARTITION BY
        { [LINEAR] HASH(expr)
        | [LINEAR] KEY [ALGORITHM={1|2}] (column_list) }
      [SUBPARTITIONS num]
    ]
    [(partition_definition [, partition_definition] ...)]

partition_definition:
    PARTITION partition_name
        [VALUES
            {LESS THAN {(expr | value_list) | MAXVALUE}
            |
            IN (value_list)}]
        [[STORAGE] ENGINE [=] engine_name]
        [COMMENT [=] 'comment_text' ]
        [DATA DIRECTORY [=] 'data_dir']
        [INDEX DIRECTORY [=] 'index_dir']
        [MAX_ROWS [=] max_number_of_rows]
        [MIN_ROWS [=] min_number_of_rows]
        [TABLESPACE [=] tablespace_name]
        [(subpartition_definition [, subpartition_definition] ...)]

subpartition_definition:
    SUBPARTITION logical_name
        [[STORAGE] ENGINE [=] engine_name]
        [COMMENT [=] 'comment_text' ]
        [DATA DIRECTORY [=] 'data_dir']
        [INDEX DIRECTORY [=] 'index_dir']
        [MAX_ROWS [=] max_number_of_rows]
        [MIN_ROWS [=] min_number_of_rows]
        [TABLESPACE [=] tablespace_name]

query_expression:
    SELECT ...   (Some valid select or union statement)
```

`CREATE TABLE` creates a table with the given name. You must have the `CREATE` privilege for the table.

By default, tables are created in the default database, using the `InnoDB` storage engine. An error occurs if the table exists, if there is no default database, or if the database does not exist.

For information about the physical representation of a table, see Section 13.1.15.2, “Files Created by CREATE TABLE”.

The original `CREATE TABLE` statement, including all specifications and table options are stored by MySQL when the table is created. For more information, see Section 13.1.15.1, “CREATE TABLE Statement Retention”.

There are several aspects to the `CREATE TABLE` statement, described under the following topics in this section:

- Table Name
- Temporary Tables
- Cloning or Copying a Table
- Column Data Types and Attributes
- Indexes and Foreign Keys
- Table Options
- Creating Partitioned Tables

Table Name

- *`tbl_name`*

  The table name can be specified as *`db_name.tbl_name`* to create the table in a specific database. This works regardless of whether there is a default database, assuming that the database exists. If you use quoted identifiers, quote the database and table names separately. For example, write `` `mydb`.`mytbl` ``, not `` `mydb.mytbl` ``.

  Rules for permissible table names are given in Section 9.2, “Schema Object Names”.

- `IF NOT EXISTS`

  Prevents an error from occurring if the table exists. However, there is no verification that the existing table has a structure identical to that indicated by the `CREATE TABLE` statement.

Temporary Tables

You can use the `TEMPORARY` keyword when creating a table. A `TEMPORARY` table is visible only to the current session, and is dropped automatically when the session is closed. For more information, see Section 13.1.15.3, “CREATE TEMPORARY TABLE Syntax”.

Cloning or Copying a Table

- `LIKE`

  Use `CREATE TABLE ... LIKE` to create an empty table based on the definition of another table, including any column attributes and indexes defined in the original table:

  ```
  CREATE TABLE new_tbl LIKE orig_tbl;
  ```

  For more information, see Section 13.1.15.4, “CREATE TABLE ... LIKE Syntax”.

- `[AS]` *`query_expression`*

  To create one table from another, add a `SELECT` statement at the end of the `CREATE TABLE` statement:

```
CREATE TABLE new_tbl AS SELECT * FROM orig_tbl;
```

For more information, see Section 13.1.15.5, “CREATE TABLE ... SELECT Syntax”.

- `IGNORE|REPLACE`

  The `IGNORE` and `REPLACE` options indicate how to handle rows that duplicate unique key values when copying a table using a `SELECT` statement.

  For more information, see Section 13.1.15.5, “CREATE TABLE ... SELECT Syntax”.

Column Data Types and Attributes

There is a hard limit of 4096 columns per table, but the effective maximum may be less for a given table and depends on the factors discussed in Section C.10.4, “Limits on Table Column Count and Row Size”.

- *`data_type`*

  `data_type` represents the data type in a column definition. *`spatial_type`* represents a spatial data type. The data type syntax shown is representative only. For a full description of the syntax available for specifying column data types, as well as information about the properties of each type, see Chapter 11, *Data Types*, and Section 11.5, “Extensions for Spatial Data”. A `JSON` data type is also supported for table columns; see Section 11.6, “The JSON Data Type”, for more information.

  - Some attributes do not apply to all data types. `AUTO_INCREMENT` applies only to integer and floating-point types. `DEFAULT` does not apply to the `BLOB`, `TEXT`, `GEOMETRY`, and `JSON` types.
  - Character data types (`CHAR`, `VARCHAR`, `TEXT`) can include `CHARACTER SET` and `COLLATE` attributes to specify the character set and collation for the column. For details, see Section 10.1, “Character Set Support”. `CHARSET` is a synonym for `CHARACTER SET`. Example:

    ```
    CREATE TABLE t (c CHAR(20) CHARACTER SET utf8 COLLATE utf8_bin);
    ```

    MySQL 8.0 interprets length specifications in character column definitions in characters. Lengths for `BINARY` and `VARBINARY` are in bytes.

  - For `CHAR`, `VARCHAR`, `BINARY`, and `VARBINARY` columns, indexes can be created that use only the leading part of column values, using *`col_name(length)`* syntax to specify an index prefix length. `BLOB` and `TEXT` columns also can be indexed, but a prefix length *must* be given. Prefix lengths are given in characters for nonbinary string types and in bytes for binary string types. That is, index entries consist of the first *`length`* characters of each column value for `CHAR`, `VARCHAR`, and `TEXT` columns, and the first *`length`* bytes of each column value for `BINARY`, `VARBINARY`, and `BLOB` columns. Indexing only a prefix of column values like this can make the index file much smaller. For additional information about index prefixes, see Section 13.1.12, “CREATE INDEX Syntax”.

    Only the `InnoDB` and `MyISAM` storage engines support indexing on `BLOB` and `TEXT` columns. For example:

    ```
    CREATE TABLE test (blob_col BLOB, INDEX(blob_col(10)));
    ```

  - `JSON` columns cannot be indexed. You can work around this restriction by creating an index on a generated column that extracts a scalar value from the `JSON` column. See Indexing a Generated Column to Provide a JSON Column Index, for a detailed example.

- `NOT NULL | NULL`

  If neither `NULL` nor `NOT NULL` is specified, the column is treated as though `NULL` had been specified.

In MySQL 8.0, only the `InnoDB`, `MyISAM`, and `MEMORY` storage engines support indexes on columns that can have `NULL` values. In other cases, you must declare indexed columns as `NOT NULL` or an error results.

- `DEFAULT`

  Specifies a default value for a column. With one exception, the default value must be a constant; it cannot be a function or an expression. This means, for example, that you cannot set the default for a date column to be the value of a function such as `NOW()` or `CURRENT_DATE`. The exception is that you can specify `CURRENT_TIMESTAMP` as the default for a `TIMESTAMP` or `DATETIME` column. See Section 11.3.5, “Automatic Initialization and Updating for TIMESTAMP and DATETIME”.

  If a column definition includes no explicit `DEFAULT` value, MySQL determines the default value as described in Section 11.7, “Data Type Default Values”.

  `BLOB`, `TEXT`, and `JSON` columns cannot be assigned a default value.

  If the `NO_ZERO_DATE` or `NO_ZERO_IN_DATE` SQL mode is enabled and a date-valued default is not correct according to that mode, `CREATE TABLE` produces a warning if strict SQL mode is not enabled and an error if strict mode is enabled. For example, with `NO_ZERO_IN_DATE` enabled, `c1 DATE DEFAULT '2010-00-00'` produces a warning.

- `AUTO_INCREMENT`

  An integer or floating-point column can have the additional attribute `AUTO_INCREMENT`. When you insert a value of `NULL` (recommended) or `0` into an indexed `AUTO_INCREMENT` column, the column is set to the next sequence value. Typically this is *`value`*`+1`, where *`value`* is the largest value for the column currently in the table. `AUTO_INCREMENT` sequences begin with `1`.

  To retrieve an `AUTO_INCREMENT` value after inserting a row, use the `LAST_INSERT_ID()` SQL function or the `mysql_insert_id()` C API function. See Section 12.14, “Information Functions”, and Section 27.7.7.38, “mysql_insert_id()”.

  If the `NO_AUTO_VALUE_ON_ZERO` SQL mode is enabled, you can store `0` in `AUTO_INCREMENT` columns as `0` without generating a new sequence value. See Section 5.1.8, “Server SQL Modes”.

  There can be only one `AUTO_INCREMENT` column per table, it must be indexed, and it cannot have a `DEFAULT` value. An `AUTO_INCREMENT` column works properly only if it contains only positive values. Inserting a negative number is regarded as inserting a very large positive number. This is done to avoid precision problems when numbers “wrap” over from positive to negative and also to ensure that you do not accidentally get an `AUTO_INCREMENT` column that contains `0`.

  For `MyISAM` tables, you can specify an `AUTO_INCREMENT` secondary column in a multiple-column key. See Section 3.6.9, “Using AUTO_INCREMENT”.

  To make MySQL compatible with some ODBC applications, you can find the `AUTO_INCREMENT` value for the last inserted row with the following query:

  ```
  SELECT * FROM tbl_name WHERE auto_col IS NULL
  ```

  This method requires that `sql_auto_is_null` variable is not set to 0. See Section 5.1.5, “Server System Variables”.

  For information about `InnoDB` and `AUTO_INCREMENT`, see Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”. For information about `AUTO_INCREMENT` and MySQL Replication, see Section 18.4.1.1, “Replication and AUTO_INCREMENT”.

- `COMMENT`

A comment for a column can be specified with the `COMMENT` option, up to 1024 characters long. The comment is displayed by the `SHOW CREATE TABLE` and `SHOW FULL COLUMNS` statements.

- `COLUMN_FORMAT`

  Used by MySQL Cluster to determine a column's storage format. This option currently has no effect on columns of tables using storage engines other than `NDB`. In MySQL 8.0 and later, `COLUMN_FORMAT` is silently ignored.

- `GENERATED ALWAYS`

  Used to specify a generated column expression. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

  Stored generated columns can be indexed. `InnoDB` supports secondary indexes on virtual generated columns. See Section 13.1.15.9, “Secondary Indexes and Generated Columns”.

Indexes and Foreign Keys

- `CONSTRAINT` *`symbol`*

  If the `CONSTRAINT` *`symbol`* clause is given, the *`symbol`* value, if used, must be unique in the database. A duplicate *`symbol`* results in an error. If the clause is not given, or a *`symbol`* is not included following the `CONSTRAINT` keyword, a name for the constraint is created automatically.

- `PRIMARY KEY`

  A unique index where all key columns must be defined as `NOT NULL`. If they are not explicitly declared as `NOT NULL`, MySQL declares them so implicitly (and silently). A table can have only one `PRIMARY KEY`. The name of a `PRIMARY KEY` is always `PRIMARY`, which thus cannot be used as the name for any other kind of index.

  If you do not have a `PRIMARY KEY` and an application asks for the `PRIMARY KEY` in your tables, MySQL returns the first `UNIQUE` index that has no `NULL` columns as the `PRIMARY KEY`.

  In `InnoDB` tables, keep the `PRIMARY KEY` short to minimize storage overhead for secondary indexes. Each secondary index entry contains a copy of the primary key columns for the corresponding row. (See Section 15.8.2.1, “Clustered and Secondary Indexes”.)

  In the created table, a `PRIMARY KEY` is placed first, followed by all `UNIQUE` indexes, and then the nonunique indexes. This helps the MySQL optimizer to prioritize which index to use and also more quickly to detect duplicated `UNIQUE` keys.

  A `PRIMARY KEY` can be a multiple-column index. However, you cannot create a multiple-column index using the `PRIMARY KEY` key attribute in a column specification. Doing so only marks that single column as primary. You must use a separate `PRIMARY KEY(`*`index_col_name, ...`*`)` clause.

  If a `PRIMARY KEY` consists of only one column that has an integer type, you can also refer to the column as `_rowid` in `SELECT` statements.

  In MySQL, the name of a `PRIMARY KEY` is `PRIMARY`. For other indexes, if you do not assign a name, the index is assigned the same name as the first indexed column, with an optional suffix (`_2`, `_3`, `...`) to make it unique. You can see index names for a table using `SHOW INDEX FROM` *`tbl_name`*. See Section 13.7.5.22, “SHOW INDEX Syntax”.

- `KEY | INDEX`

  `KEY` is normally a synonym for `INDEX`. The key attribute `PRIMARY KEY` can also be specified as just `KEY` when given in a column definition. This was implemented for compatibility with other database systems.

- `UNIQUE`

  A `UNIQUE` index creates a constraint such that all values in the index must be distinct. An error occurs if you try to add a new row with a key value that matches an existing row. For all engines, a `UNIQUE` index permits multiple `NULL` values for columns that can contain `NULL`.

  If a `UNIQUE` index consists of only one column that has an integer type, you can also refer to the column as `_rowid` in `SELECT` statements.

- `FULLTEXT`

  A `FULLTEXT` index is a special type of index used for full-text searches. Only the `InnoDB` and `MyISAM` storage engines support `FULLTEXT` indexes. They can be created only from `CHAR`, `VARCHAR`, and `TEXT` columns. Indexing always happens over the entire column; column prefix indexing is not supported and any prefix length is ignored if specified. See Section 12.9, “Full-Text Search Functions”, for details of operation. A `WITH PARSER` clause can be specified as an *`index_option`* value to associate a parser plugin with the index if full-text indexing and searching operations need special handling. This clause is valid only for `FULLTEXT` indexes. `InnoDB` and `MyISAM` support full-text parser plugins. See Full-Text Parser Plugins and Section 28.2.4.4, “Writing Full-Text Parser Plugins” for more information.

- `SPATIAL`

  You can create `SPATIAL` indexes on spatial data types. Spatial types are supported only for `InnoDB` and `MyISAM` tables, and indexed columns must be declared as `NOT NULL`. See Section 11.5, “Extensions for Spatial Data”.

- `FOREIGN KEY`

  MySQL supports foreign keys, which let you cross-reference related data across tables, and foreign key constraints, which help keep this spread-out data consistent. For definition and option information, see *`reference_definition`*, and *`reference_option`*.

  Partitioned tables employing the `InnoDB` storage engine do not support foreign keys. See Section 22.6, “Restrictions and Limitations on Partitioning”, for more information.

- `CHECK`

  The `CHECK` clause is parsed but ignored by all storage engines. See Section 1.8.2.3, “Foreign Key Differences”.

- *`index_col_name`*

  - An *`index_col_name`* specification can end with `ASC` or `DESC` to specify whether index values are stored in ascending or descending order. The default is ascending if no order specifier is given.

  - Prefixes, defined by the *`length`* attribute, can be up to 767 bytes long for `InnoDB` tables that use the `REDUNDANT` or `COMPACT` row format. The length limit is raised to 3072 bytes for `InnoDB` tables that use the `DYNAMIC` or `COMPRESSED` row format. For MyISAM tables, the prefix limit is 1000 bytes.

    Prefix limits are measured in bytes, whereas the prefix length in `CREATE TABLE`, `ALTER TABLE`, and `CREATE INDEX` statements is interpreted as number of characters for nonbinary string types (`CHAR`, `VARCHAR`, `TEXT`) and number of bytes for binary string types (`BINARY`, `VARBINARY`, `BLOB`). Take this into account when specifying a prefix length for a nonbinary string column that uses a multibyte character set.

- *`index_type`*

  Some storage engines permit you to specify an index type when creating an index. The syntax for the *`index_type`* specifier is `USING` *`type_name`*.

Example:

```
CREATE TABLE lookup
  (id INT, INDEX USING BTREE (id))
  ENGINE = MEMORY;
```

The preferred position for `USING` is after the index column list. It can be given before the column list, but support for use of the option in that position is deprecated and will be removed in a future MySQL release.

- *`index_option`*

  `index_option` values specify additional options for an index.

  - `KEY_BLOCK_SIZE`

    For `MyISAM` tables, `KEY_BLOCK_SIZE` optionally specifies the size in bytes to use for index key blocks. The value is treated as a hint; a different size could be used if necessary. A `KEY_BLOCK_SIZE` value specified for an individual index definition overrides the table-level `KEY_BLOCK_SIZE` value.

    For information about the table-level `KEY_BLOCK_SIZE` attribute, see Table Options.

  - `WITH PARSER`

    The `WITH PARSER` option can only be used with `FULLTEXT` indexes. It associates a parser plugin with the index if full-text indexing and searching operations need special handling. `InnoDB` and `MyISAM` support full-text parser plugins. If you have a `MyISAM` table with an associated full-text parser plugin, you can convert the table to `InnoDB` using `ALTER TABLE`.

  - `COMMENT`

    In MySQL 8.0, index definitions can include an optional comment of up to 1024 characters.

    You can set the `InnoDB` `MERGE_THRESHOLD` value for an individual index using the *`index_option`* `COMMENT` clause. See Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

  For more information about permissible *`index_option`* values, see Section 13.1.12, “CREATE INDEX Syntax”. For more information about indexes, see Section 8.3.1, “How MySQL Uses Indexes”.

- *`reference_definition`*

  For *`reference_definition`* syntax details and examples, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”. For information specific to foreign keys in `InnoDB`, see Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

  `InnoDB` tables support checking of foreign key constraints. The columns of the referenced table must always be explicitly named. Both `ON DELETE` and `ON UPDATE` actions on foreign keys. For more detailed information and examples, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”. For information specific to foreign keys in `InnoDB`, see Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

  For other storage engines, MySQL Server parses and ignores the `FOREIGN KEY` and `REFERENCES` syntax in `CREATE TABLE` statements. See Section 1.8.2.3, “Foreign Key Differences”.

Important

For users familiar with the ANSI/ISO SQL Standard, please note that no storage engine, including `InnoDB`, recognizes or enforces the `MATCH` clause used in referential integrity constraint definitions. Use of an explicit `MATCH` clause will not have the specified effect, and also causes `ON DELETE` and `ON UPDATE` clauses to be ignored. For these reasons, specifying `MATCH` should be avoided.

The `MATCH` clause in the SQL standard controls how `NULL` values in a composite (multiple-column) foreign key are handled when comparing to a primary key. `InnoDB` essentially implements the semantics defined by `MATCH SIMPLE`, which permit a foreign key to be all or partially `NULL`. In that case, the (child table) row containing such a foreign key is permitted to be inserted, and does not match any row in the referenced (parent) table. It is possible to implement other semantics using triggers.

Additionally, MySQL requires that the referenced columns be indexed for performance. However, `InnoDB` does not enforce any requirement that the referenced columns be declared `UNIQUE` or `NOT NULL`. The handling of foreign key references to nonunique keys or keys that contain `NULL` values is not well defined for operations such as `UPDATE` or `DELETE CASCADE`. You are advised to use foreign keys that reference only keys that are both `UNIQUE` (or `PRIMARY`) and `NOT NULL`.

MySQL parses but ignores “inline `REFERENCES` specifications” (as defined in the SQL standard) where the references are defined as part of the column specification. MySQL accepts `REFERENCES` clauses only when specified as part of a separate `FOREIGN KEY` specification.

- *`reference_option`*

  For information about the `RESTRICT`, `CASCADE`, `SET NULL`, `NO ACTION`, and `SET DEFAULT` options, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”.

Table Options

Table options are used to optimize the behavior of the table. In most cases, you do not have to specify any of them. These options apply to all storage engines unless otherwise indicated. Options that do not apply to a given storage engine may be accepted and remembered as part of the table definition. Such options then apply if you later use `ALTER TABLE` to convert the table to use a different storage engine.

- `ENGINE`

  Specifies the storage engine for the table, using one of the names shown in the following table. The engine name can be unquoted or quoted. The quoted name `'DEFAULT'` is recognized but ignored.

Storage Engine	Description
`InnoDB`	Transaction-safe tables with row locking and foreign keys. The default storage engine for new tables. See Chapter 15, *The InnoDB Storage Engine*, and in particular Section 15.1, “Introduction to InnoDB” if you have MySQL experience but are new to `InnoDB`.
`MyISAM`	The binary portable storage engine that is primarily used for read-only or read-mostly workloads. See Section 16.2, “The MyISAM Storage Engine”.
`MEMORY`	The data for this storage engine is stored only in memory. See Section 16.3, “The MEMORY Storage Engine”.

Storage Engine	Description
CSV	Tables that store rows in comma-separated values format. See Section 16.4, “The CSV Storage Engine”.
ARCHIVE	The archiving storage engine. See Section 16.5, “The ARCHIVE Storage Engine”.
EXAMPLE	An example engine. See Section 16.9, “The EXAMPLE Storage Engine”.
FEDERATED	Storage engine that accesses remote tables. See Section 16.8, “The FEDERATED Storage Engine”.
HEAP	This is a synonym for MEMORY.
MERGE	A collection of MyISAM tables used as one table. Also known as MRG_MyISAM. See Section 16.7, “The MERGE Storage Engine”.

By default, if a storage engine is specified that is not available, the statement fails with an error. You can override this behavior by removing NO_ENGINE_SUBSTITUTION from the server SQL mode (see Section 5.1.8, “Server SQL Modes”) so that MySQL allows substitution of the specified engine with the default storage engine instead. Normally in such cases, this is InnoDB, which is the default value for the default_storage_engine system variable. When NO_ENGINE_SUBSTITUTION is disabled, a warning occurs if the storage engine specification is not honored.

- AUTO_INCREMENT

  The initial AUTO_INCREMENT value for the table. In MySQL 8.0, this works for MyISAM, MEMORY, InnoDB, and ARCHIVE tables. To set the first auto-increment value for engines that do not support the AUTO_INCREMENT table option, insert a “dummy” row with a value one less than the desired value after creating the table, and then delete the dummy row.

  For engines that support the AUTO_INCREMENT table option in CREATE TABLE statements, you can also use ALTER TABLE *tbl_name* AUTO_INCREMENT = *N* to reset the AUTO_INCREMENT value. The value cannot be set lower than the maximum value currently in the column.

- AVG_ROW_LENGTH

  An approximation of the average row length for your table. You need to set this only for large tables with variable-size rows.

  When you create a MyISAM table, MySQL uses the product of the MAX_ROWS and AVG_ROW_LENGTH options to decide how big the resulting table is. If you don't specify either option, the maximum size for MyISAM data and index files is 256TB by default. (If your operating system does not support files that large, table sizes are constrained by the file size limit.) If you want to keep down the pointer sizes to make the index smaller and faster and you don't really need big files, you can decrease the default pointer size by setting the myisam_data_pointer_size system variable. (See Section 5.1.5, “Server System Variables”.) If you want all your tables to be able to grow above the default limit and are willing to have your tables slightly slower and larger than necessary, you can increase the default pointer size by setting this variable. Setting the value to 7 permits table sizes up to 65,536TB.

- [DEFAULT] CHARACTER SET

  Specifies a default character set for the table. CHARSET is a synonym for CHARACTER SET. If the character set name is DEFAULT, the database character set is used.

- CHECKSUM

  Set this to 1 if you want MySQL to maintain a live checksum for all rows (that is, a checksum that MySQL updates automatically as the table changes). This makes the table a little slower to update, but also makes it easier to find corrupted tables. The CHECKSUM TABLE statement reports the checksum. (MyISAM only.)

- `[DEFAULT] COLLATE`

  Specifies a default collation for the table.

- `COMMENT`

  A comment for the table, up to 2048 characters long.

  You can set the `InnoDB` `MERGE_THRESHOLD` value for a table using the *`table_option`* `COMMENT` clause. See Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

- `COMPRESSION`

  The compression algorithm used for page level compression for `InnoDB` tables. Supported values include `Zlib`, `LZ4`, and `None`. The `COMPRESSION` attribute was introduced with the transparent page compression feature. Page compression is only supported with `InnoDB` tables that reside in file-per-table tablespaces, and is only available on Linux and Windows platforms that support sparse files and hole punching. For more information, see Section 15.9.2, “InnoDB Page Compression”.

- `CONNECTION`

  The connection string for a `FEDERATED` table.

  Note

  Older versions of MySQL used a `COMMENT` option for the connection string.

- `DATA DIRECTORY`, `INDEX DIRECTORY`

  For `InnoDB`, the `DATA DIRECTORY='`*`directory`*`'` option allows you to create `InnoDB` file-per-table tablespaces outside the MySQL data directory. Within the directory that you specify, MySQL creates a subdirectory corresponding to the database name, and within that a `.ibd` file for the table. The `innodb_file_per_table` configuration option must be enabled to use the `DATA DIRECTORY` option with `InnoDB`. The full directory path must be specified. See Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory” for more information.

  When creating `MyISAM` tables, you can use the `DATA DIRECTORY='`*`directory`*`'` clause, the `INDEX DIRECTORY='`*`directory`*`'` clause, or both. They specify where to put a `MyISAM` table's data file and index file, respectively. Unlike `InnoDB` tables, MySQL does not create subdirectories that correspond to the database name when creating a `MyISAM` table with a `DATA DIRECTORY` or `INDEX DIRECTORY` option. Files are created in the directory that is specified.

  You must have the `FILE` privilege to use the `DATA DIRECTORY` or `INDEX DIRECTORY` table option.

  Important

  Table-level `DATA DIRECTORY` and `INDEX DIRECTORY` options are ignored for partitioned tables. (Bug #32091)

  These options work only when you are not using the `--skip-symbolic-links` option. Your operating system must also have a working, thread-safe `realpath()` call. See Section 8.12.2.2, “Using Symbolic Links for MyISAM Tables on Unix”, for more complete information.

  If a `MyISAM` table is created with no `DATA DIRECTORY` option, the `.MYD` file is created in the database directory. By default, if `MyISAM` finds an existing `.MYD` file in this case, it overwrites it. The same applies to `.MYI` files for tables created with no `INDEX DIRECTORY` option. To suppress this behavior, start the server with the `--keep_files_on_create` option, in which case `MyISAM` will not overwrite existing files and returns an error instead.

If a MyISAM table is created with a DATA DIRECTORY or INDEX DIRECTORY option and an existing .MYD or .MYI file is found, MyISAM always returns an error. It will not overwrite a file in the specified directory.

Important

You cannot use path names that contain the MySQL data directory with DATA DIRECTORY or INDEX DIRECTORY. This includes partitioned tables and individual table partitions. (See Bug #32167.)

- DELAY_KEY_WRITE

  Set this to 1 if you want to delay key updates for the table until the table is closed. See the description of the delay_key_write system variable in Section 5.1.5, “Server System Variables”. (MyISAM only.)

- ENCRYPTION

  Set the ENCRYPTION option to 'Y' to enable page-level data encryption for an InnoDB table created in a file-per-table tablespace. Option values are not case sensitive. The ENCRYPTION option was introduced with the InnoDB tablespace encryption feature; see Section 15.7.10, “InnoDB Tablespace Encryption”. The keyring_file plugin must be loaded to use the ENCRYPTION option.

- INSERT_METHOD

  If you want to insert data into a MERGE table, you must specify with INSERT_METHOD the table into which the row should be inserted. INSERT_METHOD is an option useful for MERGE tables only. Use a value of FIRST or LAST to have inserts go to the first or last table, or a value of NO to prevent inserts. See Section 16.7, “The MERGE Storage Engine”.

- KEY_BLOCK_SIZE

  For MyISAM tables, KEY_BLOCK_SIZE optionally specifies the size in bytes to use for index key blocks. The value is treated as a hint; a different size could be used if necessary. A KEY_BLOCK_SIZE value specified for an individual index definition overrides the table-level KEY_BLOCK_SIZE value.

  For InnoDB tables, KEY_BLOCK_SIZE optionally specifies the page size (in kilobytes) to use for compressed InnoDB tables. The KEY_BLOCK_SIZE value is treated as a hint; a different size could be used by InnoDB if necessary. KEY_BLOCK_SIZE can only be less than or equal to the innodb_page_size value. A value of 0 represents the default compressed page size, which is half of the innodb_page_size value. Depending on innodb_page_size, possible KEY_BLOCK_SIZE values include 0, 1, 2, 4, 8, and 16. See Section 15.9.1, “InnoDB Table Compression” for more information.

  Oracle recommends enabling innodb_strict_mode when specifying KEY_BLOCK_SIZE for InnoDB tables. When innodb_strict_mode is enabled, specifying an invalid KEY_BLOCK_SIZE value returns an error. If innodb_strict_mode is disabled, an invalid KEY_BLOCK_SIZE value results in a warning, and the KEY_BLOCK_SIZE option is ignored.

  InnoDB only supports KEY_BLOCK_SIZE at the table level.

  KEY_BLOCK_SIZE is not supported with 32k and 64k innodb_page_size values. InnoDB table compression does not support these pages sizes.

  InnoDB does not support the KEY_BLOCK_SIZE option when creating temporary tables.

- MAX_ROWS

The maximum number of rows you plan to store in the table. This is not a hard limit, but rather a hint to the storage engine that the table must be able to store at least this many rows.

The maximum `MAX_ROWS` value is 4294967295; larger values are truncated to this limit.

- `MIN_ROWS`

  The minimum number of rows you plan to store in the table. The `MEMORY` storage engine uses this option as a hint about memory use.

- `PACK_KEYS`

  Takes effect only with `MyISAM` tables. Set this option to 1 if you want to have smaller indexes. This usually makes updates slower and reads faster. Setting the option to 0 disables all packing of keys. Setting it to `DEFAULT` tells the storage engine to pack only long `CHAR`, `VARCHAR`, `BINARY`, or `VARBINARY` columns.

  If you do not use `PACK_KEYS`, the default is to pack strings, but not numbers. If you use `PACK_KEYS=1`, numbers are packed as well.

  When packing binary number keys, MySQL uses prefix compression:

  - Every key needs one extra byte to indicate how many bytes of the previous key are the same for the next key.
  - The pointer to the row is stored in high-byte-first order directly after the key, to improve compression.

  This means that if you have many equal keys on two consecutive rows, all following “same” keys usually only take two bytes (including the pointer to the row). Compare this to the ordinary case where the following keys takes `storage_size_for_key + pointer_size` (where the pointer size is usually 4). Conversely, you get a significant benefit from prefix compression only if you have many numbers that are the same. If all keys are totally different, you use one byte more per key, if the key is not a key that can have `NULL` values. (In this case, the packed key length is stored in the same byte that is used to mark if a key is `NULL`.)

- `PASSWORD`

  This option is unused.

- `ROW_FORMAT`

  Defines the physical format in which the rows are stored.

  When executing a `CREATE TABLE` statement, if you specify a row format that is not supported by the storage engine that is used for the table, the table is created using that storage engine's default row format. The information reported in this column in response to `SHOW TABLE STATUS` is the actual row format used. This may differ from the value in the `Create_options` column because the original `CREATE TABLE` definition is retained during creation.

  Row format choices differ depending on the storage engine used for the table.

  For `InnoDB` tables:

  - The default row format is defined by `innodb_default_row_format`, which has a default setting of `DYNAMIC`. The default row format is used when the `ROW_FORMAT` option is not defined or when `ROW_FORMAT=DEFAULT` is used.

    If the `ROW_FORMAT` option is not defined, or if `ROW_FORMAT=DEFAULT` is used, operations that rebuild a table also silently change the row format of the table to the default defined by

innodb_default_row_format. For more information, see Section 15.10.2, “Specifying the Row Format for a Table”.

- For more efficient InnoDB storage of data types, especially BLOB types, use the DYNAMIC. See Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats” for requirements associated with the DYNAMIC row format.
- To enable compression for InnoDB tables, specify ROW_FORMAT=COMPRESSED. The ROW_FORMAT=COMPRESSED option is not supported when creating temporary tables. See Section 15.9, “InnoDB Table and Page Compression” for requirements associated with the COMPRESSED row format.
- The row format used in older versions of MySQL can still be requested by specifying the REDUNDANT row format.
- When you specify a non-default ROW_FORMAT clause, consider also enabling the innodb_strict_mode configuration option.
- ROW_FORMAT=FIXED is not supported. If ROW_FORMAT=FIXED is specified while innodb_strict_mode is disabled, InnoDB issues a warning and assumes ROW_FORMAT=DYNAMIC. If ROW_FORMAT=FIXED is specified while innodb_strict_mode is enabled, which is the default, InnoDB returns an error.
- For additional information about InnoDB row formats, see Section 15.10, “InnoDB Row Storage and Row Formats”.

For MyISAM tables, the option value can be FIXED or DYNAMIC for static or variable-length row format. myisampack sets the type to COMPRESSED. See Section 16.2.3, “MyISAM Table Storage Formats”.

- STATS_AUTO_RECALC

  Specifies whether to automatically recalculate persistent statistics for an InnoDB table. The value DEFAULT causes the persistent statistics setting for the table to be determined by the innodb_stats_auto_recalc configuration option. The value 1 causes statistics to be recalculated when 10% of the data in the table has changed. The value 0 prevents automatic recalculation for this table; with this setting, issue an ANALYZE TABLE statement to recalculate the statistics after making substantial changes to the table. For more information about the persistent statistics feature, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- STATS_PERSISTENT

  Specifies whether to enable persistent statistics for an InnoDB table. The value DEFAULT causes the persistent statistics setting for the table to be determined by the innodb_stats_persistent configuration option. The value 1 enables persistent statistics for the table, while the value 0 turns off this feature. After enabling persistent statistics through a CREATE TABLE or ALTER TABLE statement, issue an ANALYZE TABLE statement to calculate the statistics, after loading representative data into the table. For more information about the persistent statistics feature, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- STATS_SAMPLE_PAGES

  The number of index pages to sample when estimating cardinality and other statistics for an indexed column, such as those calculated by ANALYZE TABLE. For more information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- TABLESPACE

  The TABLESPACE option may be used to create a table in an existing general tablespace, a file-per-table tablespace, or the system tablespace.

```
CREATE TABLE tbl_name ... TABLESPACE [=] tablespace_name
```

For information about general tablespaces, see Section 15.7.9, “InnoDB General Tablespaces”.

The *`tablespace_name`* is a case-sensitive identifier. It may be quoted or unquoted. The forward slash character (“/”) is not permitted. Names beginning with “innodb_” are reserved for special use.

The `TABLESPACE` option may be used to assign `InnoDB` table partitions or subpartitions to a general tablespace, a separate file-per-table tablespace, or the system tablespace. All partitions must belong to the same storage engine.

A tablespace specified at the table level becomes the default tablespace for new partitions and subpartitions. The default tablespace may be overridden by specifying a tablespace at the partition or subpartition level in a `CREATE TABLE` or `ALTER TABLE` statement. The following example shows tablespaces defined at the table level and partition level.

```
mysql> CREATE TABLE t1 ( a INT NOT NULL, PRIMARY KEY (a))
    -> ENGINE=InnoDB TABLESPACE ts1
    -> PARTITION BY RANGE (a) PARTITIONS 3 (
    -> PARTITION P1 VALUES LESS THAN (2),
    -> PARTITION P2 VALUES LESS THAN (4) TABLESPACE ts2,
    -> PARTITION P3 VALUES LESS THAN (6) TABLESPACE ts3);
```

For more information about the `TABLESPACE` option and partitioning, see Section 15.7.9, “InnoDB General Tablespaces”

To create a table in the system tablespace, specify `innodb_system` as the tablespace name.

```
CREATE TABLE tbl_name ... TABLESPACE [=] innodb_system
```

Using the `TABLESPACE [=] innodb_system` option, you can place a table of any uncompressed row format in the system tablespace regardless of the `innodb_file_per_table` setting. For example, you can add a table with `ROW_FORMAT=DYNAMIC` to the system tablespace using the `TABLESPACE [=] innodb_system` option.

To create a table in a file-per-table tablespace, specify `innodb_file_per_table` as the tablespace name.

```
CREATE TABLE tbl_name ... TABLESPACE [=] innodb_file_per_table
```

Note

If `innodb_file_per_table` is enabled, you need not specify `TABLESPACE=innodb_file_per_table` to create an `InnoDB` file-per-table tablespace. `InnoDB` tables are created in file-per-table tablespaces by default when `innodb_file_per_table` is enabled.

The `DATA DIRECTORY` clause is permitted with `CREATE TABLE ... TABLESPACE=innodb_file_per_table` but is otherwise not supported for use in combination with the `TABLESPACE` option.

The `TABLESPACE` option is supported with `ALTER TABLE` and `ALTER TABLE ... REORGANIZE PARTITION` statements, which can be used to move tables and partitions from one tablespace to another, respectively. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

- `UNION`

  Used to access a collection of identical `MyISAM` tables as one. This works only with `MERGE` tables. See Section 16.7, “The MERGE Storage Engine”.

  You must have `SELECT`, `UPDATE`, and `DELETE` privileges for the tables you map to a `MERGE` table.

Note

Formerly, all tables used had to be in the same database as the `MERGE` table itself. This restriction no longer applies.

Creating Partitioned Tables

`partition_options` can be used to control partitioning of the table created with `CREATE TABLE`.

Not all options shown in the syntax for *`partition_options`* at the beginning of this section are available for all partitioning types. Please see the listings for the following individual types for information specific to each type, and see Chapter 22, *Partitioning*, for more complete information about the workings of and uses for partitioning in MySQL, as well as additional examples of table creation and other statements relating to MySQL partitioning.

Partitions can be modified, merged, added to tables, and dropped from tables. For basic information about the MySQL statements to accomplish these tasks, see Section 13.1.7, “ALTER TABLE Syntax”. For more detailed descriptions and examples, see Section 22.3, “Partition Management”.

- `PARTITION BY`

  If used, a *`partition_options`* clause begins with `PARTITION BY`. This clause contains the function that is used to determine the partition; the function returns an integer value ranging from 1 to *`num`*, where *`num`* is the number of partitions. (The maximum number of user-defined partitions which a table may contain is 1024; the number of subpartitions—discussed later in this section—is included in this maximum.)

  Note

  The expression (*`expr`*) used in a `PARTITION BY` clause cannot refer to any columns not in the table being created; such references are specifically not permitted and cause the statement to fail with an error. (Bug #29444)

- `HASH(`*`expr`*`)`

  Hashes one or more columns to create a key for placing and locating rows. *`expr`* is an expression using one or more table columns. This can be any valid MySQL expression (including MySQL functions) that yields a single integer value. For example, these are both valid `CREATE TABLE` statements using `PARTITION BY HASH`:

  ```
  CREATE TABLE t1 (col1 INT, col2 CHAR(5))
      PARTITION BY HASH(col1);

  CREATE TABLE t1 (col1 INT, col2 CHAR(5), col3 DATETIME)
      PARTITION BY HASH ( YEAR(col3) );
  ```

  You may not use either `VALUES LESS THAN` or `VALUES IN` clauses with `PARTITION BY HASH`.

  `PARTITION BY HASH` uses the remainder of *`expr`* divided by the number of partitions (that is, the modulus). For examples and additional information, see Section 22.2.4, “HASH Partitioning”.

  The `LINEAR` keyword entails a somewhat different algorithm. In this case, the number of the partition in which a row is stored is calculated as the result of one or more logical `AND` operations. For discussion and examples of linear hashing, see Section 22.2.4.1, “LINEAR HASH Partitioning”.

- `KEY(`*`column_list`*`)`

  This is similar to `HASH`, except that MySQL supplies the hashing function so as to guarantee an even data distribution. The *`column_list`* argument is simply a list of 1 or more table columns (maximum: 16). This example shows a simple table partitioned by key, with 4 partitions:

```
CREATE TABLE tk (col1 INT, col2 CHAR(5), col3 DATE)
    PARTITION BY KEY(col3)
    PARTITIONS 4;
```

For tables that are partitioned by key, you can employ linear partitioning by using the `LINEAR` keyword. This has the same effect as with tables that are partitioned by `HASH`. That is, the partition number is found using the `&` operator rather than the modulus (see Section 22.2.4.1, “LINEAR HASH Partitioning”, and Section 22.2.5, “KEY Partitioning”, for details). This example uses linear partitioning by key to distribute data between 5 partitions:

```
CREATE TABLE tk (col1 INT, col2 CHAR(5), col3 DATE)
    PARTITION BY LINEAR KEY(col3)
    PARTITIONS 5;
```

The `ALGORITHM={1|2}` option is supported with `[SUB]PARTITION BY [LINEAR] KEY`. `ALGORITHM=1` causes the server to use the same key-hashing functions as MySQL 5.1; `ALGORITHM=2` means that the server employs the key-hashing functions implemented and used by default for new `KEY` partitioned tables in MySQL 5.5 and later. (Partitioned tables created with the key-hashing functions employed in MySQL 5.5 and later cannot be used by a MySQL 5.1 server.) Not specifying the option has the same effect as using `ALGORITHM=2`. This option is intended for use chiefly when upgrading or downgrading `[LINEAR] KEY` partitioned tables between MySQL 5.1 and later MySQL versions, or for creating tables partitioned by `KEY` or `LINEAR KEY` on a MySQL 5.5 or later server which can be used on a MySQL 5.1 server. For more information, see Section 13.1.7.1, “ALTER TABLE Partition Operations”.

`mysqldump` in MySQL 5.7 (and later) writes this option encased in versioned comments, like this:

```
CREATE TABLE t1 (a INT)
/*!50100 PARTITION BY KEY */ /*!50611 ALGORITHM = 1 */ /*!50100 ()
      PARTITIONS 3 */
```

This causes MySQL 5.6.10 and earlier servers to ignore the option, which would otherwise cause a syntax error in those versions. If you plan to load a dump made on a MySQL 5.7 server where you use tables that are partitioned or subpartitioned by `KEY` into a MySQL 5.6 server previous to version 5.6.11, be sure to consult Changes Affecting Upgrades to MySQL 5.6, before proceeding. (The information found there also applies if you are loading a dump containing `KEY` partitioned or subpartitioned tables made from a MySQL 5.7—actually 5.6.11 or later—server into a MySQL 5.5.30 or earlier server.)

Also in MySQL 5.6.11 and later, `ALGORITHM=1` is shown when necessary in the output of `SHOW CREATE TABLE` using versioned comments in the same manner as `mysqldump`. `ALGORITHM=2` is always omitted from `SHOW CREATE TABLE` output, even if this option was specified when creating the original table.

You may not use either `VALUES LESS THAN` or `VALUES IN` clauses with `PARTITION BY KEY`.

- `RANGE(`*`expr`*`)`

  In this case, *`expr`* shows a range of values using a set of `VALUES LESS THAN` operators. When using range partitioning, you must define at least one partition using `VALUES LESS THAN`. You cannot use `VALUES IN` with range partitioning.

Note

For tables partitioned by `RANGE`, `VALUES LESS THAN` must be used with either an integer literal value or an expression that evaluates to a single integer value. In MySQL 8.0, you can overcome this limitation in a table that

is defined using `PARTITION BY RANGE COLUMNS`, as described later in this section.

Suppose that you have a table that you wish to partition on a column containing year values, according to the following scheme.

Partition Number:	Years Range:
0	1990 and earlier
1	1991 to 1994
2	1995 to 1998
3	1999 to 2002
4	2003 to 2005
5	2006 and later

A table implementing such a partitioning scheme can be realized by the `CREATE TABLE` statement shown here:

```
CREATE TABLE t1 (
    year_col  INT,
    some_data INT
)
PARTITION BY RANGE (year_col) (
    PARTITION p0 VALUES LESS THAN (1991),
    PARTITION p1 VALUES LESS THAN (1995),
    PARTITION p2 VALUES LESS THAN (1999),
    PARTITION p3 VALUES LESS THAN (2002),
    PARTITION p4 VALUES LESS THAN (2006),
    PARTITION p5 VALUES LESS THAN MAXVALUE
);
```

`PARTITION ... VALUES LESS THAN ...` statements work in a consecutive fashion. `VALUES LESS THAN MAXVALUE` works to specify “leftover” values that are greater than the maximum value otherwise specified.

`VALUES LESS THAN` clauses work sequentially in a manner similar to that of the `case` portions of a `switch ... case` block (as found in many programming languages such as C, Java, and PHP). That is, the clauses must be arranged in such a way that the upper limit specified in each successive `VALUES LESS THAN` is greater than that of the previous one, with the one referencing `MAXVALUE` coming last of all in the list.

- `RANGE COLUMNS(`*`column_list`*`)`

  This variant on `RANGE` facilitates partition pruning for queries using range conditions on multiple columns (that is, having conditions such as `WHERE a = 1 AND b < 10` or `WHERE a = 1 AND b = 10 AND c < 10`). It enables you to specify value ranges in multiple columns by using a list of columns in the `COLUMNS` clause and a set of column values in each `PARTITION ... VALUES LESS THAN (`*`value_list`*`)` partition definition clause. (In the simplest case, this set consists of a single column.) The maximum number of columns that can be referenced in the *`column_list`* and *`value_list`* is 16.

  The *`column_list`* used in the `COLUMNS` clause may contain only names of columns; each column in the list must be one of the following MySQL data types: the integer types; the string types; and time or date column types. Columns using `BLOB`, `TEXT`, `SET`, `ENUM`, `BIT`, or spatial data types are not permitted; columns that use floating-point number types are also not permitted. You also may not use functions or arithmetic expressions in the `COLUMNS` clause.

  The `VALUES LESS THAN` clause used in a partition definition must specify a literal value for each column that appears in the `COLUMNS()` clause; that is, the list of values used for each `VALUES`

`LESS THAN` clause must contain the same number of values as there are columns listed in the `COLUMNS` clause. An attempt to use more or fewer values in a `VALUES LESS THAN` clause than there are in the `COLUMNS` clause causes the statement to fail with the error `Inconsistency in usage of column lists for partitioning...`. You cannot use `NULL` for any value appearing in `VALUES LESS THAN`. It is possible to use `MAXVALUE` more than once for a given column other than the first, as shown in this example:

```
CREATE TABLE rc (
    a INT NOT NULL,
    b INT NOT NULL
)
PARTITION BY RANGE COLUMNS(a,b) (
    PARTITION p0 VALUES LESS THAN (10,5),
    PARTITION p1 VALUES LESS THAN (20,10),
    PARTITION p2 VALUES LESS THAN (50,MAXVALUE),
    PARTITION p3 VALUES LESS THAN (65,MAXVALUE),
    PARTITION p4 VALUES LESS THAN (MAXVALUE,MAXVALUE)
);
```

Each value used in a `VALUES LESS THAN` value list must match the type of the corresponding column exactly; no conversion is made. For example, you cannot use the string `'1'` for a value that matches a column that uses an integer type (you must use the numeral `1` instead), nor can you use the numeral `1` for a value that matches a column that uses a string type (in such a case, you must use a quoted string: `'1'`).

For more information, see Section 22.2.1, “RANGE Partitioning”, and Section 22.4, “Partition Pruning”.

- `LIST(`*`expr`*`)`

  This is useful when assigning partitions based on a table column with a restricted set of possible values, such as a state or country code. In such a case, all rows pertaining to a certain state or country can be assigned to a single partition, or a partition can be reserved for a certain set of states or countries. It is similar to `RANGE`, except that only `VALUES IN` may be used to specify permissible values for each partition.

  `VALUES IN` is used with a list of values to be matched. For instance, you could create a partitioning scheme such as the following:

```
CREATE TABLE client_firms (
    id   INT,
    name VARCHAR(35)
)
PARTITION BY LIST (id) (
    PARTITION r0 VALUES IN (1, 5, 9, 13, 17, 21),
    PARTITION r1 VALUES IN (2, 6, 10, 14, 18, 22),
    PARTITION r2 VALUES IN (3, 7, 11, 15, 19, 23),
    PARTITION r3 VALUES IN (4, 8, 12, 16, 20, 24)
);
```

  When using list partitioning, you must define at least one partition using `VALUES IN`. You cannot use `VALUES LESS THAN` with `PARTITION BY LIST`.

  Note

  For tables partitioned by `LIST`, the value list used with `VALUES IN` must consist of integer values only. In MySQL 8.0, you can overcome this limitation using partitioning by `LIST COLUMNS`, which is described later in this section.

- `LIST COLUMNS(`*`column_list`*`)`

  This variant on `LIST` facilitates partition pruning for queries using comparison conditions on multiple columns (that is, having conditions such as `WHERE a = 5 AND b = 5` or `WHERE a = 1 AND`

`b = 10 AND c = 5`). It enables you to specify values in multiple columns by using a list of columns in the `COLUMNS` clause and a set of column values in each `PARTITION ... VALUES IN (value_list)` partition definition clause.

The rules governing regarding data types for the column list used in `LIST COLUMNS(column_list)` and the value list used in `VALUES IN(value_list)` are the same as those for the column list used in `RANGE COLUMNS(column_list)` and the value list used in `VALUES LESS THAN(value_list)`, respectively, except that in the `VALUES IN` clause, `MAXVALUE` is not permitted, and you may use `NULL`.

There is one important difference between the list of values used for `VALUES IN` with `PARTITION BY LIST COLUMNS` as opposed to when it is used with `PARTITION BY LIST`. When used with `PARTITION BY LIST COLUMNS`, each element in the `VALUES IN` clause must be a *set* of column values; the number of values in each set must be the same as the number of columns used in the `COLUMNS` clause, and the data types of these values must match those of the columns (and occur in the same order). In the simplest case, the set consists of a single column. The maximum number of columns that can be used in the *`column_list`* and in the elements making up the *`value_list`* is 16.

The table defined by the following `CREATE TABLE` statement provides an example of a table using `LIST COLUMNS` partitioning:

```
CREATE TABLE lc (
    a INT NULL,
    b INT NULL
)
PARTITION BY LIST COLUMNS(a,b) (
    PARTITION p0 VALUES IN( (0,0), (NULL,NULL) ),
    PARTITION p1 VALUES IN( (0,1), (0,2), (0,3), (1,1), (1,2) ),
    PARTITION p2 VALUES IN( (1,0), (2,0), (2,1), (3,0), (3,1) ),
    PARTITION p3 VALUES IN( (1,3), (2,2), (2,3), (3,2), (3,3) )
);
```

- `PARTITIONS` *`num`*

  The number of partitions may optionally be specified with a `PARTITIONS` *`num`* clause, where *`num`* is the number of partitions. If both this clause *and* any `PARTITION` clauses are used, *`num`* must be equal to the total number of any partitions that are declared using `PARTITION` clauses.

  Note

  Whether or not you use a `PARTITIONS` clause in creating a table that is partitioned by `RANGE` or `LIST`, you must still include at least one `PARTITION VALUES` clause in the table definition (see below).

- `SUBPARTITION BY`

  A partition may optionally be divided into a number of subpartitions. This can be indicated by using the optional `SUBPARTITION BY` clause. Subpartitioning may be done by `HASH` or `KEY`. Either of these may be `LINEAR`. These work in the same way as previously described for the equivalent partitioning types. (It is not possible to subpartition by `LIST` or `RANGE`.)

  The number of subpartitions can be indicated using the `SUBPARTITIONS` keyword followed by an integer value.

- Rigorous checking of the value used in `PARTITIONS` or `SUBPARTITIONS` clauses is applied and this value must adhere to the following rules:

  - The value must be a positive, nonzero integer.
  - No leading zeros are permitted.

- The value must be an integer literal, and cannot not be an expression. For example, `PARTITIONS 0.2E+01` is not permitted, even though `0.2E+01` evaluates to `2`. (Bug #15890)

- *`partition_definition`*

  Each partition may be individually defined using a *`partition_definition`* clause. The individual parts making up this clause are as follows:

  - `PARTITION` *`partition_name`*

    Specifies a logical name for the partition.

  - `VALUES`

    For range partitioning, each partition must include a `VALUES LESS THAN` clause; for list partitioning, you must specify a `VALUES IN` clause for each partition. This is used to determine which rows are to be stored in this partition. See the discussions of partitioning types in Chapter 22, *Partitioning*, for syntax examples.

  - `[STORAGE] ENGINE`

    MySQL accepts a `[STORAGE] ENGINE` option for both `PARTITION` and `SUBPARTITION`. Currently, the only way in which this option can be used is to set all partitions or all subpartitions to the same storage engine, and an attempt to set different storage engines for partitions or subpartitions in the same table will give rise to the error `ERROR 1469 (HY000): The mix of handlers in the partitions is not permitted in this version of MySQL`.

  - `COMMENT`

    An optional `COMMENT` clause may be used to specify a string that describes the partition. Example:

    ```
    COMMENT = 'Data for the years previous to 1999'
    ```

    The maximum length for a partition comment is 1024 characters.

  - `DATA DIRECTORY` and `INDEX DIRECTORY`

    `DATA DIRECTORY` and `INDEX DIRECTORY` may be used to indicate the directory where, respectively, the data and indexes for this partition are to be stored. Both the *`data_dir`* and the *`index_dir`* must be absolute system path names.

    You must have the `FILE` privilege to use the `DATA DIRECTORY` or `INDEX DIRECTORY` partition option.

    Example:

    ```
    CREATE TABLE th (id INT, name VARCHAR(30), adate DATE)
    PARTITION BY LIST(YEAR(adate))
    (
      PARTITION p1999 VALUES IN (1995, 1999, 2003)
        DATA DIRECTORY = '/var/appdata/95/data'
        INDEX DIRECTORY = '/var/appdata/95/idx',
      PARTITION p2000 VALUES IN (1996, 2000, 2004)
        DATA DIRECTORY = '/var/appdata/96/data'
        INDEX DIRECTORY = '/var/appdata/96/idx',
      PARTITION p2001 VALUES IN (1997, 2001, 2005)
        DATA DIRECTORY = '/var/appdata/97/data'
        INDEX DIRECTORY = '/var/appdata/97/idx',
      PARTITION p2002 VALUES IN (1998, 2002, 2006)
        DATA DIRECTORY = '/var/appdata/98/data'
        INDEX DIRECTORY = '/var/appdata/98/idx'
    );
    ```

`DATA DIRECTORY` and `INDEX DIRECTORY` behave in the same way as in the `CREATE TABLE` statement's *`table_option`* clause as used for `MyISAM` tables.

One data directory and one index directory may be specified per partition. If left unspecified, the data and indexes are stored by default in the table's database directory.

The `DATA DIRECTORY` and `INDEX DIRECTORY` options are ignored for creating partitioned tables if `NO_DIR_IN_CREATE` is in effect.

- `MAX_ROWS` and `MIN_ROWS`

  May be used to specify, respectively, the maximum and minimum number of rows to be stored in the partition. The values for *`max_number_of_rows`* and *`min_number_of_rows`* must be positive integers. As with the table-level options with the same names, these act only as “suggestions” to the server and are not hard limits.

- `TABLESPACE`

  May be used to assign `InnoDB` table partitions or subpartitions to a general tablespace, a separate file-per-table tablespace, or the system tablespace. All partitions must belong to the same storage engine. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

- *`subpartition_definition`*

  The partition definition may optionally contain one or more *`subpartition_definition`* clauses. Each of these consists at a minimum of the `SUBPARTITION` *`name`*, where *`name`* is an identifier for the subpartition. Except for the replacement of the `PARTITION` keyword with `SUBPARTITION`, the syntax for a subpartition definition is identical to that for a partition definition.

  Subpartitioning must be done by `HASH` or `KEY`, and can be done only on `RANGE` or `LIST` partitions. See Section 22.2.6, “Subpartitioning”.

Partitioning by Generated Columns

Partitioning by generated columns is permitted. For example:

```
CREATE TABLE t1 (
  s1 INT,
  s2 INT AS (EXP(s1)) STORED
)
PARTITION BY LIST (s2) (
  PARTITION p1 VALUES IN (1)
);
```

Partitioning sees a generated column as a regular column, which enables workarounds for limitations on functions that are not permitted for partitioning (see Section 22.6.3, “Partitioning Limitations Relating to Functions”). The preceding example demonstrates this technique: `EXP()` cannot be used directly in the `PARTITION BY` clause, but a generated column defined using `EXP()` is permitted.

13.1.15.1 CREATE TABLE Statement Retention

The original `CREATE TABLE` statement, including all specifications and table options are stored by MySQL when the table is created. The information is retained so that if you change storage engines, collations or other settings using an `ALTER TABLE` statement, the original table options specified are retained. This enables you to change between `InnoDB` and `MyISAM` table types even though the row formats supported by the two engines are different.

Because the text of the original statement is retained, but due to the way that certain values and options may be silently reconfigured (such as the `ROW_FORMAT`), the active table definition (accessible through `DESCRIBE` or with `SHOW TABLE STATUS`) and the table creation string (accessible through `SHOW CREATE TABLE`) will report different values.

13.1.15.2 Files Created by CREATE TABLE

For an `InnoDB` table created in a file-per-table tablespace or general tablespace, table data and associated indexes are stored in an ibd file in the database directory. When an `InnoDB` table is created in the system tablespace, table data and indexes are stored in the ibdata* files that represent the system tablespace. The `innodb_file_per_table` option controls whether tables are created in file-per-table tablespaces or the system tablespace, by default. The `TABLESPACE` option can be used to place a table in a file-per-table tablespace, general tablespace, or the system tablespace, regardless of the `innodb_file_per_table` setting.

For `MyISAM` tables, the storage engine creates data and index files. Thus, for each `MyISAM` table *`tbl_name`*, there are two disk files.

File	Purpose
`tbl_name``.MYD`	Data file
`tbl_name``.MYI`	Index file

Chapter 16, *Alternative Storage Engines*, describes what files each storage engine creates to represent tables. If a table name contains special characters, the names for the table files contain encoded versions of those characters as described in Section 9.2.3, “Mapping of Identifiers to File Names”.

13.1.15.3 CREATE TEMPORARY TABLE Syntax

You can use the `TEMPORARY` keyword when creating a table. A `TEMPORARY` table is visible only to the current session, and is dropped automatically when the session is closed. This means that two different sessions can use the same temporary table name without conflicting with each other or with an existing non-`TEMPORARY` table of the same name. (The existing table is hidden until the temporary table is dropped.)

`InnoDB` does not support compressed temporary tables. When `innodb_strict_mode` is enabled (the default), `CREATE TEMPORARY TABLE` returns an error if `ROW_FORMAT=COMPRESSED` or `KEY_BLOCK_SIZE` is specified. If `innodb_strict_mode` is disabled, warnings are issued and the temporary table is created using a non-compressed row format. `InnoDB` temporary tables are created in the shared temporary tablespace, `ibtmp1`. The `innodb_file_per-table` option does not affect the creation of `InnoDB` temporary tables.

`CREATE TABLE` does not automatically commit the current active transaction if you use the `TEMPORARY` keyword.

`TEMPORARY` tables have a very loose relationship with databases (schemas). Dropping a database does not automatically drop any `TEMPORARY` tables created within that database. Also, you can create a `TEMPORARY` table in a nonexistent database if you qualify the table name with the database name in the `CREATE TABLE` statement. In this case, all subsequent references to the table must be qualified with the database name.

To create a temporary table, you must have the `CREATE TEMPORARY TABLES` privilege. After a session has created a temporary table, the server performs no further privilege checks on the table. The creating session can perform any operation on the table, such as `DROP TABLE`, `INSERT`, `UPDATE`, or `SELECT`.

One implication of this behavior is that a session can manipulate its temporary tables even if the current user has no privilege to create them. Suppose that the current user does not have the `CREATE TEMPORARY TABLES` privilege but is able to execute a `DEFINER`-context stored procedure that executes with the privileges of a user who does have `CREATE TEMPORARY TABLES` and that creates a temporary table. While the procedure executes, the session uses the privileges of the defining user. After the procedure returns, the effective privileges revert to those of the current user, which can still see the temporary table and perform any operation on it.

As of MySQL 8.0.2, you cannot use `CREATE TEMPORY TABLE ... LIKE` to create an empty table based on the definition of a table that resides in the `mysql` tablespace, `InnoDB` system tablespace (`innodb_system`), or a general tablespace. The tablespace definition for such a table includes a `TABLESPACE` attribute that defines the tablespace where the table resides, and the aforementioned tablespaces do not support temporary tables. To create a temporary table based on the definition of such a table, use this syntax instead:

```
CREATE TEMPORARY TABLE new_tbl SELECT * FROM orig_tbl LIMIT 0;
```

13.1.15.4 CREATE TABLE ... LIKE Syntax

Use `CREATE TABLE ... LIKE` to create an empty table based on the definition of another table, including any column attributes and indexes defined in the original table:

```
CREATE TABLE new_tbl LIKE orig_tbl;
```

The copy is created using the same version of the table storage format as the original table. The `SELECT` privilege is required on the original table.

`LIKE` works only for base tables, not for views.

Important

You cannot execute `CREATE TABLE` or `CREATE TABLE ... LIKE` while a `LOCK TABLES` statement is in effect.

`CREATE TABLE ... LIKE` makes the same checks as `CREATE TABLE`. This means that if the current SQL mode is different from the mode in effect when the original table was created, the table definition might be considered invalid for the new mode and the statement will fail.

For `CREATE TABLE ... LIKE`, the destination table preserves generated column information from the original table.

`CREATE TABLE ... LIKE` does not preserve any `DATA DIRECTORY` or `INDEX DIRECTORY` table options that were specified for the original table, or any foreign key definitions.

If the original table is a `TEMPORARY` table, `CREATE TABLE ... LIKE` does not preserve `TEMPORARY`. To create a `TEMPORARY` destination table, use `CREATE TEMPORARY TABLE ... LIKE`.

As of MySQL 8.0.2, tables created in the `mysql` tablespace, the `InnoDB` system tablespace (`innodb_system`), or general tablespaces include a `TABLESPACE` attribute in the table definition, which defines the tablespace where the table resides. Due to a temporary regression introduced by this change, `CREATE TABLE ... LIKE` preserves the `TABLESPACE` attribute and creates the table in the defined tablespace regardless of the `innodb_file_per_table` setting. To avoid the `TABLESPACE` attribute when creating an empty table based on the definition of such a table, use this syntax instead:

```
CREATE TABLE new_tbl SELECT * FROM orig_tbl LIMIT 0;
```

13.1.15.5 CREATE TABLE ... SELECT Syntax

You can create one table from another by adding a `SELECT` statement at the end of the `CREATE TABLE` statement:

```
CREATE TABLE new_tbl [AS] SELECT * FROM orig_tbl;
```

MySQL creates new columns for all elements in the `SELECT`. For example:

```
mysql> CREATE TABLE test (a INT NOT NULL AUTO_INCREMENT,
    ->        PRIMARY KEY (a), KEY(b))
    ->        ENGINE=MyISAM SELECT b,c FROM test2;
```

This creates a `MyISAM` table with three columns, `a`, `b`, and `c`. The `ENGINE` option is part of the `CREATE TABLE` statement, and should not be used following the `SELECT`; this would result in a syntax error. The same is true for other `CREATE TABLE` options such as `CHARSET`.

Notice that the columns from the `SELECT` statement are appended to the right side of the table, not overlapped onto it. Take the following example:

```
mysql> SELECT * FROM foo;
+---+
| n |
+---+
| 1 |
+---+

mysql> CREATE TABLE bar (m INT) SELECT n FROM foo;
Query OK, 1 row affected (0.02 sec)
Records: 1  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM bar;
+------+---+
| m    | n |
+------+---+
| NULL | 1 |
+------+---+
1 row in set (0.00 sec)
```

For each row in table `foo`, a row is inserted in `bar` with the values from `foo` and default values for the new columns.

In a table resulting from `CREATE TABLE ... SELECT`, columns named only in the `CREATE TABLE` part come first. Columns named in both parts or only in the `SELECT` part come after that. The data type of `SELECT` columns can be overridden by also specifying the column in the `CREATE TABLE` part.

If any errors occur while copying the data to the table, it is automatically dropped and not created.

You can precede the `SELECT` by `IGNORE` or `REPLACE` to indicate how to handle rows that duplicate unique key values. With `IGNORE`, rows that duplicate an existing row on a unique key value are discarded. With `REPLACE`, new rows replace rows that have the same unique key value. If neither `IGNORE` nor `REPLACE` is specified, duplicate unique key values result in an error. For more information, see Comparison of the IGNORE Keyword and Strict SQL Mode.

Because the ordering of the rows in the underlying `SELECT` statements cannot always be determined, `CREATE TABLE ... IGNORE SELECT` and `CREATE TABLE ... REPLACE SELECT` statements are flagged as unsafe for statement-based replication. With this change, such statements produce a warning in the log when using statement-based mode and are logged using the row-based format when using `MIXED` mode. See also Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

`CREATE TABLE ... SELECT` does not automatically create any indexes for you. This is done intentionally to make the statement as flexible as possible. If you want to have indexes in the created table, you should specify these before the `SELECT` statement:

```
mysql> CREATE TABLE bar (UNIQUE (n)) SELECT n FROM foo;
```

For `CREATE TABLE ... SELECT`, the destination table does not preserve information about whether columns in the selected-from table are generated columns. The `SELECT` part of the statement cannot assign values to generated columns in the destination table.

Some conversion of data types might occur. For example, the `AUTO_INCREMENT` attribute is not preserved, and `VARCHAR` columns can become `CHAR` columns. Retrained attributes are `NULL` (or `NOT NULL`) and, for those columns that have them, `CHARACTER SET`, `COLLATION`, `COMMENT`, and the `DEFAULT` clause.

When creating a table with `CREATE TABLE ... SELECT`, make sure to alias any function calls or expressions in the query. If you do not, the `CREATE` statement might fail or result in undesirable column names.

```
CREATE TABLE artists_and_works
  SELECT artist.name, COUNT(work.artist_id) AS number_of_works
  FROM artist LEFT JOIN work ON artist.id = work.artist_id
  GROUP BY artist.id;
```

You can also explicitly specify the data type for a column in the created table:

```
CREATE TABLE foo (a TINYINT NOT NULL) SELECT b+1 AS a FROM bar;
```

For `CREATE TABLE ... SELECT`, if `IF NOT EXISTS` is given and the target table exists, nothing is inserted into the destination table, and the statement is not logged.

To ensure that the binary log can be used to re-create the original tables, MySQL does not permit concurrent inserts during `CREATE TABLE ... SELECT`.

You cannot use `FOR UPDATE` as part of the `SELECT` in a statement such as `CREATE TABLE new_table SELECT ... FROM old_table ...`. If you attempt to do so, the statement fails.

13.1.15.6 Using FOREIGN KEY Constraints

MySQL supports foreign keys, which let you cross-reference related data across tables, and foreign key constraints, which help keep this spread-out data consistent. The essential syntax for a foreign key constraint definition in a `CREATE TABLE` or `ALTER TABLE` statement looks like this:

```
[CONSTRAINT [symbol]] FOREIGN KEY
    [index_name] (index_col_name, ...)
    REFERENCES tbl_name (index_col_name,...)
    [ON DELETE reference_option]
    [ON UPDATE reference_option]

reference_option:
    RESTRICT | CASCADE | SET NULL | NO ACTION | SET DEFAULT
```

`index_name` represents a foreign key ID. The *`index_name`* value is ignored if there is already an explicitly defined index on the child table that can support the foreign key. Otherwise, MySQL implicitly creates a foreign key index that is named according to the following rules:

- If defined, the `CONSTRAINT` *`symbol`* value is used. Otherwise, the `FOREIGN KEY` *`index_name`* value is used.
- If neither a `CONSTRAINT` *`symbol`* or `FOREIGN KEY` *`index_name`* is defined, the foreign key index name is generated using the name of the referencing foreign key column.

Foreign keys definitions are subject to the following conditions:

- Foreign key relationships involve a parent table that holds the central data values, and a child table with identical values pointing back to its parent. The `FOREIGN KEY` clause is specified in the child table. The parent and child tables must use the same storage engine. They must not be `TEMPORARY` tables.

  In MySQL 8.0, creation of a foreign key constraint requires the `REFERENCES` privilege for the parent table.

- Corresponding columns in the foreign key and the referenced key must have similar data types. *The size and sign of integer types must be the same*. The length of string types need not be the same. For nonbinary (character) string columns, the character set and collation must be the same.
- When `foreign_key_checks` is enabled, which is the default setting, character set conversion is not permitted on tables that include a character string column used in a foreign key constraint. The workaround is described in Section 13.1.7, "ALTER TABLE Syntax".
- MySQL requires indexes on foreign keys and referenced keys so that foreign key checks can be fast and not require a table scan. In the referencing table, there must be an index where the foreign key columns are listed as the *first* columns in the same order. Such an index is created on the referencing table automatically if it does not exist. This index might be silently dropped later, if you create another index that can be used to enforce the foreign key constraint. *`index_name`*, if given, is used as described previously.
- `InnoDB` permits a foreign key to reference any column or group of columns. However, in the referenced table, there must be an index where the referenced columns are listed as the *first* columns in the same order.
- Index prefixes on foreign key columns are not supported. One consequence of this is that `BLOB` and `TEXT` columns cannot be included in a foreign key because indexes on those columns must always include a prefix length.
- If the `CONSTRAINT` *`symbol`* clause is given, the *`symbol`* value, if used, must be unique in the database. A duplicate *`symbol`* will result in an error similar to: `ERROR 1022 (2300): Can't write; duplicate key in table '#sql- 464_1'`. If the clause is not given, or a *`symbol`* is not included following the `CONSTRAINT` keyword, a name for the constraint is created automatically.
- `InnoDB` does not currently support foreign keys for tables with user-defined partitioning. This includes both parent and child tables.

Referential Actions

This section describes how foreign keys help guarantee referential integrity.

For storage engines supporting foreign keys, MySQL rejects any `INSERT` or `UPDATE` operation that attempts to create a foreign key value in a child table if there is no a matching candidate key value in the parent table.

When an `UPDATE` or `DELETE` operation affects a key value in the parent table that has matching rows in the child table, the result depends on the *referential action* specified using `ON UPDATE` and `ON DELETE` subclauses of the `FOREIGN KEY` clause. MySQL supports five options regarding the action to be taken, listed here:

- `CASCADE`: Delete or update the row from the parent table, and automatically delete or update the matching rows in the child table. Both `ON DELETE CASCADE` and `ON UPDATE CASCADE` are supported. Between two tables, do not define several `ON UPDATE CASCADE` clauses that act on the same column in the parent table or in the child table.

  Note

  Cascaded foreign key actions do not activate triggers.

- `SET NULL`: Delete or update the row from the parent table, and set the foreign key column or columns in the child table to `NULL`. Both `ON DELETE SET NULL` and `ON UPDATE SET NULL` clauses are supported.

  If you specify a `SET NULL` action, *make sure that you have not declared the columns in the child table as `NOT NULL`*.

- `RESTRICT`: Rejects the delete or update operation for the parent table. Specifying `RESTRICT` (or `NO ACTION`) is the same as omitting the `ON DELETE` or `ON UPDATE` clause.

- `NO ACTION`: A keyword from standard SQL. In MySQL, equivalent to `RESTRICT`. The MySQL Server rejects the delete or update operation for the parent table if there is a related foreign key value in the referenced table. Some database systems have deferred checks, and `NO ACTION` is a deferred check. In MySQL, foreign key constraints are checked immediately, so `NO ACTION` is the same as `RESTRICT`.

- `SET DEFAULT`: This action is recognized by the MySQL parser, but `InnoDB` rejects table definitions containing `ON DELETE SET DEFAULT` or `ON UPDATE SET DEFAULT` clauses.

For an `ON DELETE` or `ON UPDATE` that is not specified, the default action is always `RESTRICT`.

MySQL supports foreign key references between one column and another within a table. (A column cannot have a foreign key reference to itself.) In these cases, “child table records” really refers to dependent records within the same table.

A foreign key constraint on a stored generated column cannot use `ON UPDATE CASCADE`, `ON DELETE SET NULL`, `ON UPDATE SET NULL`, `ON DELETE SET DEFAULT`, or `ON UPDATE SET DEFAULT`.

A foreign key constraint cannot reference a virtual generated column.

For `InnoDB` restrictions related to foreign keys and generated columns, see Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

Examples of Foreign Key Clauses

Here is a simple example that relates `parent` and `child` tables through a single-column foreign key:

```
CREATE TABLE parent (
    id INT NOT NULL,
    PRIMARY KEY (id)
) ENGINE=INNODB;

CREATE TABLE child (
    id INT,
    parent_id INT,
    INDEX par_ind (parent_id),
    FOREIGN KEY (parent_id)
        REFERENCES parent(id)
        ON DELETE CASCADE
) ENGINE=INNODB;
```

A more complex example in which a `product_order` table has foreign keys for two other tables. One foreign key references a two-column index in the `product` table. The other references a single-column index in the `customer` table:

```
CREATE TABLE product (
    category INT NOT NULL, id INT NOT NULL,
    price DECIMAL,
    PRIMARY KEY(category, id)
)   ENGINE=INNODB;

CREATE TABLE customer (
    id INT NOT NULL,
    PRIMARY KEY (id)
)   ENGINE=INNODB;

CREATE TABLE product_order (
    no INT NOT NULL AUTO_INCREMENT,
    product_category INT NOT NULL,
    product_id INT NOT NULL,
    customer_id INT NOT NULL,

    PRIMARY KEY(no),
    INDEX (product_category, product_id),
```

```
    INDEX (customer_id),

    FOREIGN KEY (product_category, product_id)
      REFERENCES product(category, id)
      ON UPDATE CASCADE ON DELETE RESTRICT,

    FOREIGN KEY (customer_id)
      REFERENCES customer(id)
)   ENGINE=INNODB;
```

Adding foreign keys

You can add a new foreign key constraint to an existing table by using `ALTER TABLE`. The syntax relating to foreign keys for this statement is shown here:

```
ALTER TABLE tbl_name
    ADD [CONSTRAINT [symbol]] FOREIGN KEY
    [index_name] (index_col_name, ...)
    REFERENCES tbl_name (index_col_name,...)
    [ON DELETE reference_option]
    [ON UPDATE reference_option]
```

The foreign key can be self referential (referring to the same table). When you add a foreign key constraint to a table using `ALTER TABLE`, *remember to create the required indexes first.*

Dropping Foreign Keys

You can also use `ALTER TABLE` to drop foreign keys, using the syntax shown here:

```
ALTER TABLE tbl_name DROP FOREIGN KEY fk_symbol;
```

If the `FOREIGN KEY` clause included a `CONSTRAINT` name when you created the foreign key, you can refer to that name to drop the foreign key. Otherwise, the *`fk_symbol`* value is generated internally when the foreign key is created. To find out the symbol value when you want to drop a foreign key, use a `SHOW CREATE TABLE` statement, as shown here:

```
mysql> SHOW CREATE TABLE ibtest11c\G
*************************** 1. row ***************************
       Table: ibtest11c
Create Table: CREATE TABLE `ibtest11c` (
  `A` int(11) NOT NULL auto_increment,
  `D` int(11) NOT NULL default '0',
  `B` varchar(200) NOT NULL default '',
  `C` varchar(175) default NULL,
  PRIMARY KEY  (`A`,`D`,`B`),
  KEY `B` (`B`,`C`),
  KEY `C` (`C`),
  CONSTRAINT `0_38775` FOREIGN KEY (`A`, `D`)
REFERENCES `ibtest11a` (`A`, `D`)
ON DELETE CASCADE ON UPDATE CASCADE,
  CONSTRAINT `0_38776` FOREIGN KEY (`B`, `C`)
REFERENCES `ibtest11a` (`B`, `C`)
ON DELETE CASCADE ON UPDATE CASCADE
) ENGINE=INNODB CHARSET=utf8mb4
1 row in set (0.01 sec)

mysql> ALTER TABLE ibtest11c DROP FOREIGN KEY `0_38775`;
```

Adding and dropping a foreign key in the same `ALTER TABLE` statement is supported for `ALTER TABLE ... ALGORITHM=INPLACE` but is unsupported for `ALTER TABLE ... ALGORITHM=COPY`.

In MySQL 8.0, the server prohibits changes to foreign key columns with the potential to cause loss of referential integrity. A workaround is to use `ALTER TABLE ... DROP FOREIGN KEY` before changing the column definition and `ALTER TABLE ... ADD FOREIGN KEY` afterward.

Foreign Keys and Other MySQL Statements

Table and column identifiers in a `FOREIGN KEY ... REFERENCES ...` clause can be quoted within backticks (`). Alternatively, double quotation marks (") can be used if the `ANSI_QUOTES` SQL mode is enabled. The setting of the `lower_case_table_names` system variable is also taken into account.

You can view a child table's foreign key definitions as part of the output of the `SHOW CREATE TABLE` statement:

```
SHOW CREATE TABLE tbl_name;
```

You can also obtain information about foreign keys by querying the `INFORMATION_SCHEMA.KEY_COLUMN_USAGE` table.

You can find information about foreign keys used by `InnoDB` tables in the `INNODB_SYS_FOREIGN` and `INNODB_SYS_FOREIGN_COLS` tables, also in the `INFORMATION_SCHEMA` database.

`mysqldump` produces correct definitions of tables in the dump file, including the foreign keys for child tables.

To make it easier to reload dump files for tables that have foreign key relationships, `mysqldump` automatically includes a statement in the dump output to set `foreign_key_checks` to 0. This avoids problems with tables having to be reloaded in a particular order when the dump is reloaded. It is also possible to set this variable manually:

```
mysql> SET foreign_key_checks = 0;
mysql> SOURCE dump_file_name;
mysql> SET foreign_key_checks = 1;
```

This enables you to import the tables in any order if the dump file contains tables that are not correctly ordered for foreign keys. It also speeds up the import operation. Setting `foreign_key_checks` to 0 can also be useful for ignoring foreign key constraints during `LOAD DATA` and `ALTER TABLE` operations. However, even if `foreign_key_checks = 0`, MySQL does not permit the creation of a foreign key constraint where a column references a nonmatching column type. Also, if a table has foreign key constraints, `ALTER TABLE` cannot be used to alter the table to use another storage engine. To change the storage engine, you must drop any foreign key constraints first.

You cannot issue `DROP TABLE` for a table that is referenced by a `FOREIGN KEY` constraint, unless you do `SET foreign_key_checks = 0`. When you drop a table, any constraints that were defined in the statement used to create that table are also dropped.

If you re-create a table that was dropped, it must have a definition that conforms to the foreign key constraints referencing it. It must have the correct column names and types, and it must have indexes on the referenced keys, as stated earlier. If these are not satisfied, MySQL returns Error 1005 and refers to Error 150 in the error message, which means that a foreign key constraint was not correctly formed. Similarly, if an `ALTER TABLE` fails due to Error 150, this means that a foreign key definition would be incorrectly formed for the altered table.

For `InnoDB` tables, you can obtain a detailed explanation of the most recent `InnoDB` foreign key error in the MySQL Server, by checking the output of `SHOW ENGINE INNODB STATUS`.

Important

For users familiar with the ANSI/ISO SQL Standard, please note that no storage engine, including `InnoDB`, recognizes or enforces the `MATCH` clause used in referential-integrity constraint definitions. Use of an explicit `MATCH` clause will not have the specified effect, and also causes `ON DELETE` and `ON UPDATE` clauses to be ignored. For these reasons, specifying `MATCH` should be avoided.

The `MATCH` clause in the SQL standard controls how `NULL` values in a composite (multiple-column) foreign key are handled when comparing to a

> primary key. MySQL essentially implements the semantics defined by `MATCH SIMPLE`, which permit a foreign key to be all or partially `NULL`. In that case, the (child table) row containing such a foreign key is permitted to be inserted, and does not match any row in the referenced (parent) table. It is possible to implement other semantics using triggers.
>
> Additionally, MySQL requires that the referenced columns be indexed for performance reasons. However, the system does not enforce a requirement that the referenced columns be `UNIQUE` or be declared `NOT NULL`. The handling of foreign key references to nonunique keys or keys that contain `NULL` values is not well defined for operations such as `UPDATE` or `DELETE CASCADE`. You are advised to use foreign keys that reference only `UNIQUE` (including `PRIMARY`) and `NOT NULL` keys.
>
> Furthermore, MySQL parses but ignores “inline `REFERENCES` specifications” (as defined in the SQL standard) where the references are defined as part of the column specification. MySQL accepts `REFERENCES` clauses only when specified as part of a separate `FOREIGN KEY` specification. For storage engines that do not support foreign keys (such as `MyISAM`), MySQL Server parses and ignores foreign key specifications.

13.1.15.7 Silent Column Specification Changes

In some cases, MySQL silently changes column specifications from those given in a `CREATE TABLE` or `ALTER TABLE` statement. These might be changes to a data type, to attributes associated with a data type, or to an index specification.

All changes are subject to the internal row-size limit of 65,535 bytes, which may cause some attempts at data type changes to fail. See Section C.10.4, “Limits on Table Column Count and Row Size”.

- Columns that are part of a `PRIMARY KEY` are made `NOT NULL` even if not declared that way.
- Trailing spaces are automatically deleted from `ENUM` and `SET` member values when the table is created.
- MySQL maps certain data types used by other SQL database vendors to MySQL types. See Section 11.10, “Using Data Types from Other Database Engines”.
- If you include a `USING` clause to specify an index type that is not permitted for a given storage engine, but there is another index type available that the engine can use without affecting query results, the engine uses the available type.
- If strict SQL mode is not enabled, a `VARCHAR` column with a length specification greater than 65535 is converted to `TEXT`, and a `VARBINARY` column with a length specification greater than 65535 is converted to `BLOB`. Otherwise, an error occurs in either of these cases.
- Specifying the `CHARACTER SET binary` attribute for a character data type causes the column to be created as the corresponding binary data type: `CHAR` becomes `BINARY`, `VARCHAR` becomes `VARBINARY`, and `TEXT` becomes `BLOB`. For the `ENUM` and `SET` data types, this does not occur; they are created as declared. Suppose that you specify a table using this definition:

```
CREATE TABLE t
(
  c1 VARCHAR(10) CHARACTER SET binary,
  c2 TEXT CHARACTER SET binary,
  c3 ENUM('a','b','c') CHARACTER SET binary
);
```

The resulting table has this definition:

```
CREATE TABLE t
(
  c1 VARBINARY(10),
  c2 BLOB,
  c3 ENUM('a','b','c') CHARACTER SET binary
);
```

To see whether MySQL used a data type other than the one you specified, issue a `DESCRIBE` or `SHOW CREATE TABLE` statement after creating or altering the table.

Certain other data type changes can occur if you compress a table using `myisampack`. See Section 16.2.3.3, “Compressed Table Characteristics”.

13.1.15.8 CREATE TABLE and Generated Columns

`CREATE TABLE` supports the specification of generated columns. Values of a generated column are computed from an expression included in the column definition.

The following simple example shows a table that stores the lengths of the sides of right triangles in the `sidea` and `sideb` columns, and computes the length of the hypotenuse in `sidec` (the square root of the sums of the squares of the other sides):

```
CREATE TABLE triangle (
  sidea DOUBLE,
  sideb DOUBLE,
  sidec DOUBLE AS (SQRT(sidea * sidea + sideb * sideb))
);
INSERT INTO triangle (sidea, sideb) VALUES(1,1),(3,4),(6,8);
```

Selecting from the table yields this result:

```
mysql> SELECT * FROM triangle;
+-------+-------+--------------------+
| sidea | sideb | sidec              |
+-------+-------+--------------------+
|     1 |     1 | 1.4142135623730951 |
|     3 |     4 |                  5 |
|     6 |     8 |                 10 |
+-------+-------+--------------------+
```

Any application that uses the `triangle` table has access to the hypotenuse values without having to specify the expression that calculates them.

Generated column definitions have this syntax:

```
col_name data_type [GENERATED ALWAYS] AS (expression)
  [VIRTUAL | STORED] [UNIQUE [KEY]] [COMMENT comment]
  [[NOT] NULL] [[PRIMARY] KEY]
```

`AS (expression)` indicates that the column is generated and defines the expression used to compute column values. `AS` may be preceded by `GENERATED ALWAYS` to make the generated nature of the column more explicit. Constructs that are permitted or prohibited in the expression are discussed later.

The `VIRTUAL` or `STORED` keyword indicates how column values are stored, which has implications for column use:

- `VIRTUAL`: Column values are not stored, but are evaluated when rows are read, immediately after any `BEFORE` triggers. A virtual column takes no storage.

  `InnoDB` supports secondary indexes on virtual columns. See Section 13.1.15.9, “Secondary Indexes and Generated Columns”.

- `STORED`: Column values are evaluated and stored when rows are inserted or updated. A stored column does require storage space and can be indexed.

The default is `VIRTUAL` if neither keyword is specified.

It is permitted to mix `VIRTUAL` and `STORED` columns within a table.

Other attributes may be given to indicate whether the column is indexed or can be `NULL`, or provide a comment. (Note that the order of these attributes differs from their order in nongenerated column definitions.)

Generated column expressions must adhere to the following rules. An error occurs if an expression contains disallowed constructs.

- Literals, deterministic built-in functions, and operators are permitted. A function is deterministic if, given the same data in tables, multiple invocations produce the same result, independently of the connected user. Examples of functions that fail this definition: `CONNECTION_ID()`, `CURRENT_USER()`, `NOW()`.
- Subqueries, parameters, variables, stored functions, and user-defined functions are not permitted.
- A generated column definition can refer to other generated columns, but only those occurring earlier in the table definition. A generated column definition can refer to any base (nongenerated) column in the table whether its definition occurs earlier or later.
- The `AUTO_INCREMENT` attribute cannot be used in a generated column definition.
- An `AUTO_INCREMENT` column cannot be used as a base column in a generated column definition.
- If expression evaluation causes truncation or provides incorrect input to a function, the `CREATE TABLE` statement terminates with an error and the DDL operation is rejected.

If the expression evaluates to a data type that differs from the declared column type, coercion to the declared type occurs according to the usual MySQL type-conversion rules. See Section 12.2, “Type Conversion in Expression Evaluation”.

Note

If any component of the expression depends on the SQL mode, different results may occur for different uses of the table unless the SQL mode is the same during all uses.

For `CREATE TABLE ... LIKE`, the destination table preserves generated column information from the original table.

For `CREATE TABLE ... SELECT`, the destination table does not preserve information about whether columns in the selected-from table are generated columns. The `SELECT` part of the statement cannot assign values to generated columns in the destination table.

Partitioning by generated columns is permitted. See Creating Partitioned Tables.

A foreign key constraint on a stored generated column cannot use `ON UPDATE CASCADE`, `ON DELETE SET NULL`, `ON UPDATE SET NULL`, `ON DELETE SET DEFAULT`, or `ON UPDATE SET DEFAULT`.

A foreign key constraint cannot reference a virtual generated column.

For `InnoDB` restrictions related to foreign keys and generated columns, see Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints”.

Triggers cannot use `NEW.col_name` or use `OLD.col_name` to refer to generated columns.

For INSERT, REPLACE, and UPDATE, if a generated column is inserted into, replaced, or updated explicitly, the only permitted value is DEFAULT.

A generated column in a view is considered updatable because it is possible to assign to it. However, if such a column is updated explicitly, the only permitted value is DEFAULT.

Generated columns have several use cases, such as these:

- Virtual generated columns can be used as a way to simplify and unify queries. A complicated condition can be defined as a generated column and referred to from multiple queries on the table to ensure that all of them use exactly the same condition.
- Stored generated columns can be used as a materialized cache for complicated conditions that are costly to calculate on the fly.
- Generated columns can simulate functional indexes: Use a stored column to define a functional expression and index it. This can be useful for working with columns of types that cannot be indexed directly, such as JSON columns; see Indexing a Generated Column to Provide a JSON Column Index, for a detailed example.

  The disadvantage of such an approach is that values are stored twice; once as the value of the generated column and once in the index.
- If a generated column is indexed, the optimizer recognizes query expressions that match the column definition and uses indexes from the column as appropriate during query execution, even if a query does not refer to the column directly by name. For details, see Section 8.3.10, “Optimizer Use of Generated Column Indexes”.

Example:

Suppose that a table t1 contains first_name and last_name columns and that applications frequently construct the full name using an expression like this:

```
SELECT CONCAT(first_name,' ',last_name) AS full_name FROM t1;
```

One way to avoid writing out the expression is to create a view v1 on t1, which simplifies applications by enabling them to select full_name directly without using an expression:

```
CREATE VIEW v1 AS
SELECT *, CONCAT(first_name,' ',last_name) AS full_name FROM t1;

SELECT full_name FROM v1;
```

A generated column also enables applications to select full_name directly without the need to define a view:

```
CREATE TABLE t1 (
  first_name VARCHAR(10),
  last_name VARCHAR(10),
  full_name VARCHAR(255) AS (CONCAT(first_name,' ',last_name))
);

SELECT full_name FROM t1;
```

13.1.15.9 Secondary Indexes and Generated Columns

InnoDB supports secondary indexes on virtual generated columns. Other index types are not supported. A secondary index defined on a virtual column is sometimes referred to as a “virtual index”.

A secondary index may be created on one or more virtual columns or on a combination of virtual columns and regular columns or stored generated columns. Secondary indexes that include virtual columns may be defined as UNIQUE.

When a secondary index is created on a virtual generated column, generated column values are materialized in the records of the index. If the index is a covering index (one that includes all the columns retrieved by a query), generated column values are retrieved from materialized values in the index structure instead of computed “on the fly”.

There are additional write costs to consider when using a secondary index on a virtual column due to computation performed when materializing virtual column values in secondary index records during `INSERT` and `UPDATE` operations. Even with additional write costs, secondary indexes on virtual columns may be preferable to generated *stored* columns, which are materialized in the clustered index, resulting in larger tables that require more disk space and memory. If a secondary index is not defined on a virtual column, there are additional costs for reads, as virtual column values must be computed each time the column's row is examined.

Values of an indexed virtual column are MVCC-logged to avoid unnecessary recomputation of generated column values during rollback or during a purge operation. The data length of logged values is limited by the index key limit of 767 bytes for `COMPACT` and `REDUNDANT` row formats, and 3072 bytes for `DYNAMIC` and `COMPRESSED` row formats.

Adding or dropping a secondary index on a virtual column is an in-place operation.

Indexing a Generated Column to Provide a JSON Column Index

As noted elsewhere, `JSON` columns cannot be indexed directly. To create an index that references such a column indirectly, you can define a generated column that extracts the information that should be indexed, then create an index on the generated column, as shown in this example:

```
mysql> CREATE TABLE jemp (
    ->     c JSON,
    ->     g INT GENERATED ALWAYS AS (c->"$.id")),
    ->     INDEX i (g)
    -> );
Query OK, 0 rows affected (0.28 sec)

mysql> INSERT INTO jemp (c) VALUES
     >   ('{"id": "1", "name": "Fred"}'), ('{"id": "2", "name": "Wilma"}'),
     >   ('{"id": "3", "name": "Barney"}'), ('{"id": "4", "name": "Betty"}');
Query OK, 4 rows affected (0.04 sec)
Records: 4  Duplicates: 0  Warnings: 0

mysql> SELECT c->>"$.name" AS name
     >     FROM jemp WHERE g > 2;
+--------+
| name   |
+--------+
| Barney |
| Betty  |
+--------+
2 rows in set (0.00 sec)

mysql> EXPLAIN SELECT c->>"$.name" AS name
     >    FROM jemp WHERE g > 2\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: jemp
   partitions: NULL
         type: range
possible_keys: i
          key: i
      key_len: 5
          ref: NULL
         rows: 2
     filtered: 100.00
        Extra: Using where
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
```

```
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: /* select#1 */ select json_unquote(json_extract(`test`.`jemp`.`c`,'$.name'))
AS `name` from `test`.`jemp` where (`test`.`jemp`.`g` > 2)
1 row in set (0.00 sec)
```

(We have wrapped the output from the last statement in this example to fit the viewing area.)

When you use `EXPLAIN` on a `SELECT` or other SQL statement containing one or more expressions that use the `->` or `->>` operator, these expressions are translated into their equivalents using `JSON_EXTRACT()` and (if needed) `JSON_UNQUOTE()` instead, as shown here in the output from `SHOW WARNINGS` immediately following this `EXPLAIN` statement:

```
mysql> EXPLAIN SELECT c->>"$.name"
     > FROM jemp WHERE g > 2 ORDER BY c->"$.name"\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: jemp
   partitions: NULL
         type: range
possible_keys: i
          key: i
      key_len: 5
          ref: NULL
         rows: 2
     filtered: 100.00
        Extra: Using where; Using filesort
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Note
   Code: 1003
Message: /* select#1 */ select json_unquote(json_extract(`test`.`jemp`.`c`,'$.name')) AS
`c->>"$.name"` from `test`.`jemp` where (`test`.`jemp`.`g` > 2) order by
json_extract(`test`.`jemp`.`c`,'$.name')
1 row in set (0.00 sec)
```

See the descriptions of the `->` and `->>` operators, as well as those of the `JSON_EXTRACT()` and `JSON_UNQUOTE()` functions, for additional information and examples.

This technique also can be used to provide indexes that indirectly reference columns of other types that cannot be indexed directly, such as `GEOMETRY` columns.

13.1.16 CREATE TABLESPACE Syntax

```
CREATE TABLESPACE tablespace_name
    ADD DATAFILE 'file_name'
    [FILE_BLOCK_SIZE = value]
        [ENGINE [=] engine_name]
```

This statement is used to create an `InnoDB` tablespace. An `InnoDB` tablespace created using `CREATE TABLESPACE` is referred to as *general tablespace*.

A general tablespace is a shared tablespace, similar to the system tablespace. It can hold multiple tables, and supports all table row formats. General tablespaces can also be created in a location relative to or independent of the MySQL data directory.

After creating an `InnoDB` general tablespace, you can use `CREATE TABLE tbl_name ... TABLESPACE [=] tablespace_name` or `ALTER TABLE tbl_name TABLESPACE [=] tablespace_name` to add tables to the tablespace.

For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

Note

`CREATE TABLESPACE` is supported with `InnoDB`. In earlier releases, `CREATE TABLESPACE` only supported `NDB`, which is the MySQL NDB Cluster storage engine.

Options

- `ADD DATAFILE`: Defines the name of the tablespace data file. A data file must be specified with the `CREATE TABLESPACE` statement, and the data file name must have a `.ibd` extension. An `InnoDB` general tablespace only supports a single data file.

  To place the data file in a location outside of the MySQL data directory (`DATADIR`), include an absolute directory path or a path relative to the MySQL data directory. If you do not specify a path, the general tablespace is created in the MySQL data directory.

  To avoid conflicts with implicitly created file-per-table tablespaces, creating a general tablespace in a subdirectory under the MySQL data directory is not supported. Also, when creating a general tablespace outside of the MySQL data directory, the directory must exist prior to creating the tablespace.

  The *`file_name`*, including the path (optional), must be quoted with single or double quotations marks. File names (not counting the ".ibd" extension) and directory names must be at least one byte in length. Zero length file names and directory names are not supported.

- `FILE_BLOCK_SIZE`: Defines the block size of the tablespace data file. If you do not specify this option, `FILE_BLOCK_SIZE` defaults to `innodb_page_size`. The `FILE_BLOCK_SIZE` setting is only required if you will use the tablespace to store compressed `InnoDB` tables (`ROW_FORMAT=COMPRESSED`). In this case, you must define the tablespace `FILE_BLOCK_SIZE` when creating the tablespace.

  If `FILE_BLOCK_SIZE` is equal `innodb_page_size`, the tablespace can only contain tables with an uncompressed row format (`COMPACT`, `REDUNDANT`, and `DYNAMIC` row formats). Tables with a `COMPRESSED` row format have a different physical page size than uncompressed tables. Therefore, compressed tables cannot coexist in the same tablespace as uncompressed tables.

  For a general tablespace to contain compressed tables, `FILE_BLOCK_SIZE` must be specified, and the `FILE_BLOCK_SIZE` value must be a valid compressed page size in relation to the `innodb_page_size` value. Also, the physical page size of the compressed table (`KEY_BLOCK_SIZE`) must be equal to `FILE_BLOCK_SIZE/1024`. For example, if `innodb_page_size=16K`, and `FILE_BLOCK_SIZE=8K`, the `KEY_BLOCK_SIZE` of the table must be 8. For more information, see Section 15.7.9, "InnoDB General Tablespaces".

- `ENGINE`: Defines the storage engine which uses the tablespace, where *`engine_name`* is the name of the storage engine. Currently, only the `InnoDB` storage engine is supported. `ENGINE = InnoDB` must be defined as part of the `CREATE TABLESPACE` statement or `InnoDB` must be defined as the default storage engine (`default_storage_engine=InnoDB`).

Notes

- *`tablespace_name`* is a case-sensitive identifier for the tablespace. It may be quoted or unquoted. The forward slash character ("/") is not permitted. Names beginning with `innodb_` are either not permitted or are reserved for special use.
- Creation of temporary general tablespaces is not supported.
- General tablespaces do not support temporary tables.
- The `TABLESPACE` option may be used with `CREATE TABLE` or `ALTER TABLE` to assign `InnoDB` table partitions or subpartitions to a general tablespace, a separate file-per-table tablespace, or the

system tablespace. All partitions must belong to the same storage engine. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

- General tablespaces support the addition of tables of any row format using `CREATE TABLE ... TABLESPACE`. `innodb_file_per_table` does not need to be enabled.
- `innodb_strict_mode` is not applicable to general tablespaces. Tablespace management rules are strictly enforced independently of `innodb_strict_mode`. If `CREATE TABLESPACE` parameters are incorrect or incompatible, the operation fails regardless of the `innodb_strict_mode` setting. When a table is added to a general tablespace using `CREATE TABLE ... TABLESPACE` or `ALTER TABLE ... TABLESPACE`, `innodb_strict_mode` is ignored but the statement is evaluated as if `innodb_strict_mode` is enabled.
- Use `DROP TABLESPACE` to remove a general tablespace. All tables must be dropped from a general tablespace using `DROP TABLE` prior to dropping the tablespace.
- All parts of a table added to a general tablespace reside in the general tablespace, including indexes and `BLOB` pages.
- Similar to the system tablespace, truncating or dropping tables stored in a general tablespace creates free space internally in the general tablespace .ibd data file which can only be used for new `InnoDB` data. Space is not released back to the operating system as it is for file-per-table tablespaces.
- A general tablespace is not associated with any database or schema.
- `ALTER TABLE ... DISCARD TABLESPACE` and `ALTER TABLE ...IMPORT TABLESPACE` are not supported for tables that belong to a general tablespace.
- The server uses tablespace-level metadata locking for DDL that references general tablespaces. By comparison, the server uses table-level metadata locking for DDL that references file-per-table tablespaces.
- A generated or existing tablespace cannot be changed to a general tablespace.
- There is no conflict between general tablespace names and file-per-table tablespace names. The “/” character, which is present in file-per-table tablespace names, is not permitted in general tablespace names.

Examples

This example demonstrates creating a general tablespace and adding three uncompressed tables of different row formats.

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=REDUNDANT;
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t2 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=COMPACT;
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t3 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=DYNAMIC;
Query OK, 0 rows affected (0.00 sec)
```

This example demonstrates creating a general tablespace and adding a compressed table. The example assumes a default `innodb_page_size` of 16K. The `FILE_BLOCK_SIZE` of 8192 requires that the compressed table have a `KEY_BLOCK_SIZE` of 8.

```
mysql> CREATE TABLESPACE `ts2` ADD DATAFILE 'ts2.ibd' FILE_BLOCK_SIZE = 8192 Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)
```

```
mysql> CREATE TABLE t4 (c1 INT PRIMARY KEY) TABLESPACE ts2 ROW_FORMAT=COMPRESSED
KEY_BLOCK_SIZE=8;
Query OK, 0 rows affected (0.00 sec)
```

13.1.17 CREATE TRIGGER Syntax

```
CREATE
    [DEFINER = { user | CURRENT_USER }]
    TRIGGER trigger_name
    trigger_time trigger_event
    ON tbl_name FOR EACH ROW
    [trigger_order]
    trigger_body

trigger_time: { BEFORE | AFTER }

trigger_event: { INSERT | UPDATE | DELETE }

trigger_order: { FOLLOWS | PRECEDES } other_trigger_name
```

This statement creates a new trigger. A trigger is a named database object that is associated with a table, and that activates when a particular event occurs for the table. The trigger becomes associated with the table named *`tbl_name`*, which must refer to a permanent table. You cannot associate a trigger with a `TEMPORARY` table or a view.

Trigger names exist in the schema namespace, meaning that all triggers must have unique names within a schema. Triggers in different schemas can have the same name.

This section describes `CREATE TRIGGER` syntax. For additional discussion, see Section 23.3.1, “Trigger Syntax and Examples”.

`CREATE TRIGGER` requires the `TRIGGER` privilege for the table associated with the trigger. The statement might also require the `SET_USER_ID` or `SUPER` privilege, depending on the `DEFINER` value, as described later in this section. If binary logging is enabled, `CREATE TRIGGER` might require the `SUPER` privilege, as described in Section 23.7, “Binary Logging of Stored Programs”.

The `DEFINER` clause determines the security context to be used when checking access privileges at trigger activation time, as described later in this section.

`trigger_time` is the trigger action time. It can be `BEFORE` or `AFTER` to indicate that the trigger activates before or after each row to be modified.

Basic column value checks occur prior to trigger activation, so you cannot use `BEFORE` triggers to convert values inappropriate for the column type to valid values.

`trigger_event` indicates the kind of operation that activates the trigger. These *`trigger_event`* values are permitted:

- `INSERT`: The trigger activates whenever a new row is inserted into the table; for example, through `INSERT`, `LOAD DATA`, and `REPLACE` statements.
- `UPDATE`: The trigger activates whenever a row is modified; for example, through `UPDATE` statements.
- `DELETE`: The trigger activates whenever a row is deleted from the table; for example, through `DELETE` and `REPLACE` statements. `DROP TABLE` and `TRUNCATE TABLE` statements on the table do *not* activate this trigger, because they do not use `DELETE`. Dropping a partition does not activate `DELETE` triggers, either.

The *`trigger_event`* does not represent a literal type of SQL statement that activates the trigger so much as it represents a type of table operation. For example, an `INSERT` trigger activates not only for `INSERT` statements but also `LOAD DATA` statements because both statements insert rows into a table.

A potentially confusing example of this is the `INSERT INTO ... ON DUPLICATE KEY UPDATE ...` syntax: a `BEFORE INSERT` trigger activates for every row, followed by either an `AFTER INSERT` trigger or both the `BEFORE UPDATE` and `AFTER UPDATE` triggers, depending on whether there was a duplicate key for the row.

Note

Cascaded foreign key actions do not activate triggers.

It is possible to define multiple triggers for a given table that have the same trigger event and action time. For example, you can have two `BEFORE UPDATE` triggers for a table. By default, triggers that have the same trigger event and action time activate in the order they were created. To affect trigger order, specify a *`trigger_order`* clause that indicates `FOLLOWS` or `PRECEDES` and the name of an existing trigger that also has the same trigger event and action time. With `FOLLOWS`, the new trigger activates after the existing trigger. With `PRECEDES`, the new trigger activates before the existing trigger.

`trigger_body` is the statement to execute when the trigger activates. To execute multiple statements, use the `BEGIN ... END` compound statement construct. This also enables you to use the same statements that are permitted within stored routines. See Section 13.6.1, “BEGIN ... END Compound-Statement Syntax”. Some statements are not permitted in triggers; see Section C.1, “Restrictions on Stored Programs”.

Within the trigger body, you can refer to columns in the subject table (the table associated with the trigger) by using the aliases `OLD` and `NEW`. `OLD.col_name` refers to a column of an existing row before it is updated or deleted. `NEW.col_name` refers to the column of a new row to be inserted or an existing row after it is updated.

Triggers cannot use `NEW.col_name` or use `OLD.col_name` to refer to generated columns. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

MySQL stores the `sql_mode` system variable setting in effect when a trigger is created, and always executes the trigger body with this setting in force, *regardless of the current server SQL mode when the trigger begins executing*.

The `DEFINER` clause specifies the MySQL account to be used when checking access privileges at trigger activation time. If a *`user`* value is given, it should be a MySQL account specified as *`'user_name'@'host_name'`*, `CURRENT_USER`, or `CURRENT_USER()`. The default `DEFINER` value is the user who executes the `CREATE TRIGGER` statement. This is the same as specifying `DEFINER = CURRENT_USER` explicitly.

If you specify the `DEFINER` clause, these rules determine the valid `DEFINER` user values:

- If you do not have the `SET_USER_ID` or `SUPER` privilege, the only permitted *`user`* value is your own account, either specified literally or by using `CURRENT_USER`. You cannot set the definer to some other account.
- If you have the `SET_USER_ID` or `SUPER` privilege, you can specify any syntactically valid account name. If the account does not exist, a warning is generated.
- Although it is possible to create a trigger with a nonexistent `DEFINER` account, it is not a good idea for such triggers to be activated until the account actually does exist. Otherwise, the behavior with respect to privilege checking is undefined.

MySQL takes the `DEFINER` user into account when checking trigger privileges as follows:

- At `CREATE TRIGGER` time, the user who issues the statement must have the `TRIGGER` privilege.
- At trigger activation time, privileges are checked against the `DEFINER` user. This user must have these privileges:

- The `TRIGGER` privilege for the subject table.
- The `SELECT` privilege for the subject table if references to table columns occur using `OLD.col_name` or `NEW.col_name` in the trigger body.
- The `UPDATE` privilege for the subject table if table columns are targets of `SET NEW.col_name = value` assignments in the trigger body.
- Whatever other privileges normally are required for the statements executed by the trigger.

For more information about trigger security, see Section 23.6, “Access Control for Stored Programs and Views”.

Within a trigger body, the `CURRENT_USER()` function returns the account used to check privileges at trigger activation time. This is the `DEFINER` user, not the user whose actions caused the trigger to be activated. For information about user auditing within triggers, see Section 6.3.13, “SQL-Based MySQL Account Activity Auditing”.

If you use `LOCK TABLES` to lock a table that has triggers, the tables used within the trigger are also locked, as described in Section 13.3.5.2, “LOCK TABLES and Triggers”.

For additional discussion of trigger use, see Section 23.3.1, “Trigger Syntax and Examples”.

13.1.18 CREATE VIEW Syntax

```
CREATE
    [OR REPLACE]
    [ALGORITHM = {UNDEFINED | MERGE | TEMPTABLE}]
    [DEFINER = { user | CURRENT_USER }]
    [SQL SECURITY { DEFINER | INVOKER }]
    VIEW view_name [(column_list)]
    AS select_statement
    [WITH [CASCADED | LOCAL] CHECK OPTION]
```

The `CREATE VIEW` statement creates a new view, or replaces an existing view if the `OR REPLACE` clause is given. If the view does not exist, `CREATE OR REPLACE VIEW` is the same as `CREATE VIEW`. If the view does exist, `CREATE OR REPLACE VIEW` is the same as `ALTER VIEW`.

For information about restrictions on view use, see Section C.5, “Restrictions on Views”.

The *`select_statement`* is a `SELECT` statement that provides the definition of the view. (Selecting from the view selects, in effect, using the `SELECT` statement.) The *`select_statement`* can select from base tables or other views.

The view definition is “frozen” at creation time and is not affected by subsequent changes to the definitions of the underlying tables. For example, if a view is defined as `SELECT *` on a table, new columns added to the table later do not become part of the view, and columns dropped from the table will result in an error when selecting from the view.

The `ALGORITHM` clause affects how MySQL processes the view. The `DEFINER` and `SQL SECURITY` clauses specify the security context to be used when checking access privileges at view invocation time. The `WITH CHECK OPTION` clause can be given to constrain inserts or updates to rows in tables referenced by the view. These clauses are described later in this section.

The `CREATE VIEW` statement requires the `CREATE VIEW` privilege for the view, and some privilege for each column selected by the `SELECT` statement. For columns used elsewhere in the `SELECT` statement, you must have the `SELECT` privilege. If the `OR REPLACE` clause is present, you must also have the `DROP` privilege for the view. `CREATE VIEW` might also require the `SET_USER_ID` or `SUPER` privilege, depending on the `DEFINER` value, as described later in this section.

When a view is referenced, privilege checking occurs as described later in this section.

A view belongs to a database. By default, a new view is created in the default database. To create the view explicitly in a given database, use *db_name.view_name* syntax to qualify the view name with the database name:

```
CREATE VIEW test.v AS SELECT * FROM t;
```

Unqualified table or view names in the SELECT statement are also interpreted with respect to the default database. A view can refer to tables or views in other databases by qualifying the table or view name with the appropriate database name.

Within a database, base tables and views share the same namespace, so a base table and a view cannot have the same name.

Columns retrieved by the SELECT statement can be simple references to table columns, or expressions that use functions, constant values, operators, and so forth.

A view must have unique column names with no duplicates, just like a base table. By default, the names of the columns retrieved by the SELECT statement are used for the view column names. To define explicit names for the view columns, specify the optional *column_list* clause as a list of comma-separated identifiers. The number of names in *column_list* must be the same as the number of columns retrieved by the SELECT statement.

A view can be created from many kinds of SELECT statements. It can refer to base tables or other views. It can use joins, UNION, and subqueries. The SELECT need not even refer to any tables:

```
CREATE VIEW v_today (today) AS SELECT CURRENT_DATE;
```

The following example defines a view that selects two columns from another table as well as an expression calculated from those columns:

```
mysql> CREATE TABLE t (qty INT, price INT);
mysql> INSERT INTO t VALUES(3, 50);
mysql> CREATE VIEW v AS SELECT qty, price, qty*price AS value FROM t;
mysql> SELECT * FROM v;
+------+-------+-------+
| qty  | price | value |
+------+-------+-------+
|    3 |    50 |   150 |
+------+-------+-------+
```

A view definition is subject to the following restrictions:

- The SELECT statement cannot refer to system variables or user-defined variables.
- Within a stored program, the SELECT statement cannot refer to program parameters or local variables.
- The SELECT statement cannot refer to prepared statement parameters.
- Any table or view referred to in the definition must exist. If, after the view has been created, a table or view that the definition refers to is dropped, use of the view results in an error. To check a view definition for problems of this kind, use the CHECK TABLE statement.
- The definition cannot refer to a TEMPORARY table, and you cannot create a TEMPORARY view.
- You cannot associate a trigger with a view.
- Aliases for column names in the SELECT statement are checked against the maximum column length of 64 characters (not the maximum alias length of 256 characters).

ORDER BY is permitted in a view definition, but it is ignored if you select from a view using a statement that has its own ORDER BY.

For other options or clauses in the definition, they are added to the options or clauses of the statement that references the view, but the effect is undefined. For example, if a view definition includes a `LIMIT` clause, and you select from the view using a statement that has its own `LIMIT` clause, it is undefined which limit applies. This same principle applies to options such as `ALL`, `DISTINCT`, or `SQL_SMALL_RESULT` that follow the `SELECT` keyword, and to clauses such as `INTO`, `FOR UPDATE`, `FOR SHARE`, `LOCK IN SHARE MODE`, and `PROCEDURE`.

The results obtained from a view may be affected if you change the query processing environment by changing system variables:

```
mysql> CREATE VIEW v (mycol) AS SELECT 'abc';
Query OK, 0 rows affected (0.01 sec)

mysql> SET sql_mode = '';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT "mycol" FROM v;
+-------+
| mycol |
+-------+
| mycol |
+-------+
1 row in set (0.01 sec)

mysql> SET sql_mode = 'ANSI_QUOTES';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT "mycol" FROM v;
+-------+
| mycol |
+-------+
| abc   |
+-------+
1 row in set (0.00 sec)
```

The `DEFINER` and `SQL SECURITY` clauses determine which MySQL account to use when checking access privileges for the view when a statement is executed that references the view. The valid `SQL SECURITY` characteristic values are `DEFINER` (the default) and `INVOKER`. These indicate that the required privileges must be held by the user who defined or invoked the view, respectively.

If a *`user`* value is given for the `DEFINER` clause, it should be a MySQL account specified as *`'user_name'@'host_name'`*, `CURRENT_USER`, or `CURRENT_USER()`. The default `DEFINER` value is the user who executes the `CREATE VIEW` statement. This is the same as specifying `DEFINER = CURRENT_USER` explicitly.

If the `DEFINER` clause is present, these rules determine the valid `DEFINER` user values:

- If you do not have the `SET_USER_ID` or `SUPER` privilege, the only valid *`user`* value is your own account, either specified literally or by using `CURRENT_USER`. You cannot set the definer to some other account.
- If you have the `SET_USER_ID` or `SUPER` privilege, you can specify any syntactically valid account name. If the account does not exist, a warning is generated.
- Although it is possible to create a view with a nonexistent `DEFINER` account, an error occurs when the view is referenced if the `SQL SECURITY` value is `DEFINER` but the definer account does not exist.

For more information about view security, see Section 23.6, “Access Control for Stored Programs and Views”.

Within a view definition, `CURRENT_USER` returns the view's `DEFINER` value by default. For views defined with the `SQL SECURITY INVOKER` characteristic, `CURRENT_USER` returns the account for

the view's invoker. For information about user auditing within views, see Section 6.3.13, “SQL-Based MySQL Account Activity Auditing”.

Within a stored routine that is defined with the `SQL SECURITY DEFINER` characteristic, `CURRENT_USER` returns the routine's `DEFINER` value. This also affects a view defined within such a routine, if the view definition contains a `DEFINER` value of `CURRENT_USER`.

MySQL checks view privileges like this:

- At view definition time, the view creator must have the privileges needed to use the top-level objects accessed by the view. For example, if the view definition refers to table columns, the creator must have some privilege for each column in the select list of the definition, and the `SELECT` privilege for each column used elsewhere in the definition. If the definition refers to a stored function, only the privileges needed to invoke the function can be checked. The privileges required at function invocation time can be checked only as it executes: For different invocations, different execution paths within the function might be taken.
- The user who references a view must have appropriate privileges to access it (`SELECT` to select from it, `INSERT` to insert into it, and so forth.)
- When a view has been referenced, privileges for objects accessed by the view are checked against the privileges held by the view `DEFINER` account or invoker, depending on whether the `SQL SECURITY` characteristic is `DEFINER` or `INVOKER`, respectively.
- If reference to a view causes execution of a stored function, privilege checking for statements executed within the function depend on whether the function `SQL SECURITY` characteristic is `DEFINER` or `INVOKER`. If the security characteristic is `DEFINER`, the function runs with the privileges of the `DEFINER` account. If the characteristic is `INVOKER`, the function runs with the privileges determined by the view's `SQL SECURITY` characteristic.

Example: A view might depend on a stored function, and that function might invoke other stored routines. For example, the following view invokes a stored function `f()`:

```
CREATE VIEW v AS SELECT * FROM t WHERE t.id = f(t.name);
```

Suppose that `f()` contains a statement such as this:

```
IF name IS NULL then
  CALL p1();
ELSE
  CALL p2();
END IF;
```

The privileges required for executing statements within `f()` need to be checked when `f()` executes. This might mean that privileges are needed for `p1()` or `p2()`, depending on the execution path within `f()`. Those privileges must be checked at runtime, and the user who must possess the privileges is determined by the `SQL SECURITY` values of the view `v` and the function `f()`.

The `DEFINER` and `SQL SECURITY` clauses for views are extensions to standard SQL. In standard SQL, views are handled using the rules for `SQL SECURITY DEFINER`. The standard says that the definer of the view, which is the same as the owner of the view's schema, gets applicable privileges on the view (for example, `SELECT`) and may grant them. MySQL has no concept of a schema “owner”, so MySQL adds a clause to identify the definer. The `DEFINER` clause is an extension where the intent is to have what the standard has; that is, a permanent record of who defined the view. This is why the default `DEFINER` value is the account of the view creator.

The optional `ALGORITHM` clause is a MySQL extension to standard SQL. It affects how MySQL processes the view. `ALGORITHM` takes three values: `MERGE`, `TEMPTABLE`, or `UNDEFINED`. For more information, see Section 23.5.2, “View Processing Algorithms”, as well as Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

Some views are updatable. That is, you can use them in statements such as `UPDATE`, `DELETE`, or `INSERT` to update the contents of the underlying table. For a view to be updatable, there must be a one-to-one relationship between the rows in the view and the rows in the underlying table. There are also certain other constructs that make a view nonupdatable.

A generated column in a view is considered updatable because it is possible to assign to it. However, if such a column is updated explicitly, the only permitted value is `DEFAULT`. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

The `WITH CHECK OPTION` clause can be given for an updatable view to prevent inserts or updates to rows except those for which the `WHERE` clause in the *`select_statement`* is true.

In a `WITH CHECK OPTION` clause for an updatable view, the `LOCAL` and `CASCADED` keywords determine the scope of check testing when the view is defined in terms of another view. The `LOCAL` keyword restricts the `CHECK OPTION` only to the view being defined. `CASCADED` causes the checks for underlying views to be evaluated as well. When neither keyword is given, the default is `CASCADED`.

For more information about updatable views and the `WITH CHECK OPTION` clause, see Section 23.5.3, “Updatable and Insertable Views”, and Section 23.5.4, “The View WITH CHECK OPTION Clause”.

13.1.19 DROP DATABASE Syntax

```
DROP {DATABASE | SCHEMA} [IF EXISTS] db_name
```

`DROP DATABASE` drops all tables in the database and deletes the database. Be *very* careful with this statement! To use `DROP DATABASE`, you need the `DROP` privilege on the database. `DROP SCHEMA` is a synonym for `DROP DATABASE`.

Important

When a database is dropped, user privileges on the database are *not* automatically dropped. See Section 13.7.1.6, “GRANT Syntax”.

`IF EXISTS` is used to prevent an error from occurring if the database does not exist.

If the default database is dropped, the default database is unset (the `DATABASE()` function returns `NULL`).

If you use `DROP DATABASE` on a symbolically linked database, both the link and the original database are deleted.

`DROP DATABASE` returns the number of tables that were removed.

The `DROP DATABASE` statement removes from the given database directory those files and directories that MySQL itself may create during normal operation. This includes all files with the extensions shown in the following table.

`.BAK`	`.DAT`	`.HSH`	`.MRG`
`.MYD`	`.MYI`	`.cfg`	`.db`
`.ibd`	`.ndb`		

If other files or directories remain in the database directory after MySQL removes those just listed, the database directory cannot be removed. In this case, you must remove any remaining files or directories manually and issue the `DROP DATABASE` statement again.

Dropping a database does not remove any `TEMPORARY` tables that were created in that database. `TEMPORARY` tables are automatically removed when the session that created them ends. See Section 13.1.15.3, “CREATE TEMPORARY TABLE Syntax”.

You can also drop databases with `mysqladmin`. See Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”.

13.1.20 DROP EVENT Syntax

```
DROP EVENT [IF EXISTS] event_name
```

This statement drops the event named *`event_name`*. The event immediately ceases being active, and is deleted completely from the server.

If the event does not exist, the error `ERROR 1517 (HY000): Unknown event 'event_name'` results. You can override this and cause the statement to generate a warning for nonexistent events instead using `IF EXISTS`.

This statement requires the `EVENT` privilege for the schema to which the event to be dropped belongs.

13.1.21 DROP FUNCTION Syntax

The `DROP FUNCTION` statement is used to drop stored functions and user-defined functions (UDFs):

- For information about dropping stored functions, see Section 13.1.23, “DROP PROCEDURE and DROP FUNCTION Syntax”.
- For information about dropping user-defined functions, see Section 13.7.3.2, “DROP FUNCTION Syntax”.

13.1.22 DROP INDEX Syntax

```
DROP INDEX index_name ON tbl_name
    [algorithm_option | lock_option] ...

algorithm_option:
    ALGORITHM [=] {DEFAULT|INPLACE|COPY}

lock_option:
    LOCK [=] {DEFAULT|NONE|SHARED|EXCLUSIVE}
```

`DROP INDEX` drops the index named *`index_name`* from the table *`tbl_name`*. This statement is mapped to an `ALTER TABLE` statement to drop the index. See Section 13.1.7, “ALTER TABLE Syntax”.

To drop a primary key, the index name is always `PRIMARY`, which must be specified as a quoted identifier because `PRIMARY` is a reserved word:

```
DROP INDEX `PRIMARY` ON t;
```

`ALGORITHM` and `LOCK` clauses may be given. These influence the table copying method and level of concurrency for reading and writing the table while its indexes are being modified. They have the same meaning as for the `ALTER TABLE` statement. For more information, see Section 13.1.7, “ALTER TABLE Syntax”

13.1.23 DROP PROCEDURE and DROP FUNCTION Syntax

```
DROP {PROCEDURE | FUNCTION} [IF EXISTS] sp_name
```

This statement is used to drop a stored procedure or function. That is, the specified routine is removed from the server. You must have the `ALTER ROUTINE` privilege for the routine. (If the

`automatic_sp_privileges` system variable is enabled, that privilege and `EXECUTE` are granted automatically to the routine creator when the routine is created and dropped from the creator when the routine is dropped. See Section 23.2.2, “Stored Routines and MySQL Privileges”.)

The `IF EXISTS` clause is a MySQL extension. It prevents an error from occurring if the procedure or function does not exist. A warning is produced that can be viewed with `SHOW WARNINGS`.

`DROP FUNCTION` is also used to drop user-defined functions (see Section 13.7.3.2, “DROP FUNCTION Syntax”).

13.1.24 DROP SERVER Syntax

```
DROP SERVER [ IF EXISTS ] server_name
```

Drops the server definition for the server named *`server_name`*. The corresponding row in the `mysql.servers` table is deleted. This statement requires the `SUPER` privilege.

Dropping a server for a table does not affect any `FEDERATED` tables that used this connection information when they were created. See Section 13.1.14, “CREATE SERVER Syntax”.

`DROP SERVER` does not cause an automatic commit.

`DROP SERVER` is not written to the binary log, regardless of the logging format that is in use.

13.1.25 DROP TABLE Syntax

```
DROP [TEMPORARY] TABLE [IF EXISTS]
    tbl_name [, tbl_name] ...
    [RESTRICT | CASCADE]
```

`DROP TABLE` removes one or more tables. You must have the `DROP` privilege for each table. All table data and the table definition are *removed*, so *be careful* with this statement! If any of the tables named in the argument list do not exist, MySQL returns an error indicating by name which nonexisting tables it was unable to drop, but it also drops all of the tables in the list that do exist.

Important

When a table is dropped, user privileges on the table are *not* automatically dropped. See Section 13.7.1.6, “GRANT Syntax”.

For a partitioned table, `DROP TABLE` permanently removes the table definition, all of its partitions, and all of the data which was stored in those partitions. It also removes partition definitions associated with the dropped table.

Use `IF EXISTS` to prevent an error from occurring for tables that do not exist. A `NOTE` is generated for each nonexistent table when using `IF EXISTS`. See Section 13.7.5.40, “SHOW WARNINGS Syntax”.

`IF EXISTS` can be useful for dropping tables in unusual circumstances under which there is an entry in the data dictionary but no table managed by the storage engine. (For example, if an abnormal server exit occurs after removal of the table from the storage engine but before removal of the data dictionary entry.)

`RESTRICT` and `CASCADE` are permitted to make porting easier. In MySQL 8.0, they do nothing.

Note

`DROP TABLE` automatically commits the current active transaction, unless you use the `TEMPORARY` keyword.

The `TEMPORARY` keyword has the following effects:

- The statement drops only TEMPORARY tables.
- The statement does not end an ongoing transaction.
- No access rights are checked. (A TEMPORARY table is visible only to the session that created it, so no check is necessary.)

Using TEMPORARY is a good way to ensure that you do not accidentally drop a non-TEMPORARY table.

DROP TABLE is not supported with all innodb_force_recovery settings. See Section 15.20.2, "Forcing InnoDB Recovery".

13.1.26 DROP TABLESPACE Syntax

```
DROP TABLESPACE tablespace_name
   [ENGINE [=] engine_name]
```

This statement is used to drop an InnoDB general tablespace that was created using CREATE TABLESPACE (see Section 13.1.16, "CREATE TABLESPACE Syntax").

All tables must be dropped from the tablespace prior to a DROP TABLESPACE operation. If the tablespace is not empty, DROP TABLESPACE returns an error.

tablespace_name is a case-sensitive identifier in MySQL.

ENGINE: Defines the storage engine that uses the tablespace, where *engine_name* is the name of the storage engine. Currently, only the InnoDB storage engine is supported. You do not need to specify ENGINE = InnoDB if InnoDB is defined as the default storage engine (default_storage_engine=InnoDB).

DROP TABLESPACE is supported with InnoDB. In earlier releases, DROP TABLESPACE supports NDB, which is the MySQL Cluster storage engine. DROP TABLESPACE will support NDB in MySQL 5.7 when MySQL Cluster is branched from the MyQL 5.7 code base. The latest version of MySQL Cluster is based on MySQL 5.6.

Notes

- A general InnoDB tablespace is not deleted automatically when the last table in the tablespace is dropped. The tablespace must be dropped explicitly using DROP TABLESPACE *tablespace_name*.
- A DROP DATABASE operation can drop tables that belong to a general tablespace but it cannot drop the tablespace, even if the operation drops all tables that belong to the tablespace. The tablespace must be dropped explicitly using DROP TABLESPACE *tablespace_name*.
- Similar to the system tablespace, truncating or dropping tables stored in a general tablespace creates free space internally in the general tablespace .ibd data file which can only be used for new InnoDB data. Space is not released back to the operating system as it is for file-per-table tablespaces.

Example

This example demonstrates how to drop an InnoDB general tablespace. The general tablespace ts1 is created with a single table. Before dropping the tablespace, the table must be dropped.

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts10 Engine=InnoDB;
Query OK, 0 rows affected (0.02 sec)

mysql> DROP TABLE t1;
```

```
Query OK, 0 rows affected (0.01 sec)

mysql> DROP TABLESPACE ts1;
Query OK, 0 rows affected (0.01 sec)
```

13.1.27 DROP TRIGGER Syntax

```
DROP TRIGGER [IF EXISTS] [schema_name.]trigger_name
```

This statement drops a trigger. The schema (database) name is optional. If the schema is omitted, the trigger is dropped from the default schema. `DROP TRIGGER` requires the `TRIGGER` privilege for the table associated with the trigger.

Use `IF EXISTS` to prevent an error from occurring for a trigger that does not exist. A `NOTE` is generated for a nonexistent trigger when using `IF EXISTS`. See Section 13.7.5.40, “SHOW WARNINGS Syntax”.

Triggers for a table are also dropped if you drop the table.

13.1.28 DROP VIEW Syntax

```
DROP VIEW [IF EXISTS]
    view_name [, view_name] ...
    [RESTRICT | CASCADE]
```

`DROP VIEW` removes one or more views. You must have the `DROP` privilege for each view. If any of the views named in the argument list do not exist, MySQL returns an error indicating by name which nonexisting views it was unable to drop, but it also drops all of the views in the list that do exist.

The `IF EXISTS` clause prevents an error from occurring for views that don't exist. When this clause is given, a `NOTE` is generated for each nonexistent view. See Section 13.7.5.40, “SHOW WARNINGS Syntax”.

`RESTRICT` and `CASCADE`, if given, are parsed and ignored.

13.1.29 RENAME TABLE Syntax

```
RENAME TABLE tbl_name TO new_tbl_name
    [, tbl_name2 TO new_tbl_name2] ...
```

This statement renames one or more tables. The rename operation is done atomically, which means that no other session can access any of the tables while the rename is running.

For example, a table named `old_table` can be renamed to `new_table` as shown here:

```
RENAME TABLE old_table TO new_table;
```

This statement is equivalent to the following `ALTER TABLE` statement:

```
ALTER TABLE old_table RENAME new_table;
```

If the statement renames more than one table, renaming operations are done from left to right. If you want to swap two table names, you can do so like this (assuming that `tmp_table` does not already exist):

```
RENAME TABLE old_table TO tmp_table,
             new_table TO old_table,
```

```
            tmp_table TO new_table;
```

MySQL checks the destination table name before checking whether the source table exists. For example, if `new_table` already exists and `old_table` does not, the following statement fails as shown here:

```
mysql> SHOW TABLES;
+----------------+
| Tables_in_mydb |
+----------------+
| table_a        |
+----------------+
1 row in set (0.00 sec)

mysql> RENAME TABLE table_b TO table_a;
ERROR 1050 (42S01): Table 'table_a' already exists
```

As long as two databases are on the same file system, you can use `RENAME TABLE` to move a table from one database to another:

```
RENAME TABLE current_db.tbl_name TO other_db.tbl_name;
```

You can use this method to move all tables from one database to a different one, in effect renaming the database. (MySQL has no single statement to perform this task.)

If there are any triggers associated with a table which is moved to a different database using `RENAME TABLE`, then the statement fails with the error `Trigger in wrong schema`.

`RENAME TABLE` changes internally generated foreign key constraint names and user-defined foreign key constraint names that contain the string "*`tbl_name`*_ibfk_" to reflect the new table name. `InnoDB` interprets foreign key constraint names that contain the string "*`tbl_name`*_ibfk_" as internally generated names.

Foreign keys that point to the renamed table are not automatically updated. In such cases, you must drop and re-create the foreign keys in order for them to function properly.

`RENAME TABLE` also works for views, as long as you do not try to rename a view into a different database.

Any privileges granted specifically for the renamed table or view are not migrated to the new name. They must be changed manually.

When you execute `RENAME TABLE`, you cannot have any locked tables or active transactions. You must also have the `ALTER` and `DROP` privileges on the original table, and the `CREATE` and `INSERT` privileges on the new table.

If MySQL encounters any errors in a multiple-table rename, it does a reverse rename for all renamed tables to return everything to its original state.

You cannot use `RENAME TABLE` to rename a `TEMPORARY` table. However, you can use `ALTER TABLE` with temporary tables.

Like `RENAME TABLE`, `ALTER TABLE ... RENAME` can also be used to move a table to a different database. Regardless of the statement used to perform the rename, if the rename operation would move the table to a database located on a different file system, the success of the outcome is platform specific and depends on the underlying operating system calls used to move the table files.

13.1.30 TRUNCATE TABLE Syntax

```
TRUNCATE [TABLE] tbl_name
```

`TRUNCATE TABLE` empties a table completely. It requires the `DROP` privilege. Logically, `TRUNCATE TABLE` is similar to a `DELETE` statement that deletes all rows, or a sequence of `DROP TABLE` and `CREATE TABLE` statements.

Warning

In MySQL 8.0, `TRUNCATE TABLE` is mapped to `DROP TABLE` and `CREATE TABLE`. As a result of this change, `TRUNCATE TABLE` is temporarily non-atomic. A server exit during a `TRUNCATE TABLE` operation can result in a dropped table and orphaned foreign key constraints in the `InnoDB` `SYS_FOREIGN` and `SYS_FOREIGN_COLS` system tables.

To achieve high performance, `TRUNCATE TABLE` bypasses the DML method of deleting data. Thus, it cannot be rolled back, it does not cause `ON DELETE` triggers to fire, and it cannot be performed for `InnoDB` tables with parent-child foreign key relationships.

Although `TRUNCATE TABLE` is similar to `DELETE`, it is classified as a DDL statement rather than a DML statement. It differs from `DELETE` in the following ways:

- Truncate operations drop and re-create the table, which is much faster than deleting rows one by one, particularly for large tables.
- Truncate operations cause an implicit commit, and so cannot be rolled back.
- Truncation operations cannot be performed if the session holds an active table lock.
- `TRUNCATE TABLE` fails for an `InnoDB` table or `NDB` table if there are any `FOREIGN KEY` constraints from other tables that reference the table. Foreign key constraints between columns of the same table are permitted.
- Truncation operations do not return a meaningful value for the number of deleted rows. The usual result is “0 rows affected,” which should be interpreted as “no information.”
- As long as the table definition is valid, the table can be re-created as an empty table with `TRUNCATE TABLE`, even if the data or index files have become corrupted.
- Any `AUTO_INCREMENT` value is reset to its start value. This is true even for `MyISAM` and `InnoDB`, which normally do not reuse sequence values.
- When used with partitioned tables, `TRUNCATE TABLE` preserves the partitioning; that is, the data and index files are dropped and re-created, while the partition definitions are unaffected.
- The `TRUNCATE TABLE` statement does not invoke `ON DELETE` triggers.

`TRUNCATE TABLE` for a table closes all handlers for the table that were opened with `HANDLER OPEN`.

`TRUNCATE TABLE` is treated for purposes of binary logging and replication as `DROP TABLE` followed by `CREATE TABLE`—that is, as DDL rather than DML. This is due to the fact that, when using `InnoDB` and other transactional storage engines where the transaction isolation level does not permit statement-based logging (`READ COMMITTED` or `READ UNCOMMITTED`), the statement was not logged and replicated when using `STATEMENT` or `MIXED` logging mode. (Bug #36763) However, it is still applied on replication slaves using `InnoDB` in the manner described previously.

In MySQL 5.7 and earlier, on a system with a large buffer pool and `innodb_adaptive_hash_index` enabled, a `TRUNCATE TABLE` operation could cause a temporary drop in system performance due to an LRU scan that occurred when removing the table's adaptive hash index entries (Bug #68184). The remapping of `TRUNCATE TABLE` to `DROP TABLE` and `CREATE TABLE` in MySQL 8.0 avoids the problematic LRU scan.

`TRUNCATE TABLE` can be used with Performance Schema summary tables, but the effect is to reset the summary columns to 0 or `NULL`, not to remove rows. See Section 25.11.15, “Performance Schema Summary Tables”.

13.2 Data Manipulation Statements

13.2.1 CALL Syntax

```
CALL sp_name([parameter[,...]])
CALL sp_name[()]
```

The CALL statement invokes a stored procedure that was defined previously with CREATE PROCEDURE.

Stored procedures that take no arguments can be invoked without parentheses. That is, CALL p() and CALL p are equivalent.

CALL can pass back values to its caller using parameters that are declared as OUT or INOUT parameters. When the procedure returns, a client program can also obtain the number of rows affected for the final statement executed within the routine: At the SQL level, call the ROW_COUNT() function; from the C API, call the mysql_affected_rows() function.

To get back a value from a procedure using an OUT or INOUT parameter, pass the parameter by means of a user variable, and then check the value of the variable after the procedure returns. (If you are calling the procedure from within another stored procedure or function, you can also pass a routine parameter or local routine variable as an IN or INOUT parameter.) For an INOUT parameter, initialize its value before passing it to the procedure. The following procedure has an OUT parameter that the procedure sets to the current server version, and an INOUT value that the procedure increments by one from its current value:

```
CREATE PROCEDURE p (OUT ver_param VARCHAR(25), INOUT incr_param INT)
BEGIN
  # Set value of OUT parameter
  SELECT VERSION() INTO ver_param;
  # Increment value of INOUT parameter
  SET incr_param = incr_param + 1;
END;
```

Before calling the procedure, initialize the variable to be passed as the INOUT parameter. After calling the procedure, the values of the two variables will have been set or modified:

```
mysql> SET @increment = 10;
mysql> CALL p(@version, @increment);
mysql> SELECT @version, @increment;
+--------------+------------+
| @version     | @increment |
+--------------+------------+
| 5.5.3-m3-log |         11 |
+--------------+------------+
```

In prepared CALL statements used with PREPARE and EXECUTE, placeholders can be used for IN parameters, OUT, and INOUT parameters. These types of parameters can be used as follows:

```
mysql> SET @increment = 10;
mysql> PREPARE s FROM 'CALL p(?, ?)';
mysql> EXECUTE s USING @version, @increment;
mysql> SELECT @version, @increment;
+--------------+------------+
| @version     | @increment |
+--------------+------------+
| 5.5.3-m3-log |         11 |
+--------------+------------+
```

To write C programs that use the CALL SQL statement to execute stored procedures that produce result sets, the CLIENT_MULTI_RESULTS flag must be enabled. This is because each CALL returns

a result to indicate the call status, in addition to any result sets that might be returned by statements executed within the procedure. `CLIENT_MULTI_RESULTS` must also be enabled if `CALL` is used to execute any stored procedure that contains prepared statements. It cannot be determined when such a procedure is loaded whether those statements will produce result sets, so it is necessary to assume that they will.

`CLIENT_MULTI_RESULTS` can be enabled when you call `mysql_real_connect()`, either explicitly by passing the `CLIENT_MULTI_RESULTS` flag itself, or implicitly by passing `CLIENT_MULTI_STATEMENTS` (which also enables `CLIENT_MULTI_RESULTS`). In MySQL 8.0, `CLIENT_MULTI_RESULTS` is enabled by default.

To process the result of a `CALL` statement executed using `mysql_query()` or `mysql_real_query()`, use a loop that calls `mysql_next_result()` to determine whether there are more results. For an example, see Section 27.7.20, “C API Support for Multiple Statement Execution”.

In MySQL 8.0, C programs can use the prepared-statement interface to execute `CALL` statements and access `OUT` and `INOUT` parameters. This is done by processing the result of a `CALL` statement using a loop that calls `mysql_stmt_next_result()` to determine whether there are more results. For an example, see Section 27.7.23, “C API Support for Prepared CALL Statements”. Languages that provide a MySQL interface can use prepared `CALL` statements to directly retrieve `OUT` and `INOUT` procedure parameters.

In MySQL 8.0, metadata changes to objects referred to by stored programs are detected and cause automatic reparsing of the affected statements when the program is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

13.2.2 DELETE Syntax

`DELETE` is a DML statement that removes rows from a table.

A `DELETE` statement can start with a `WITH` clause to define common table expressions accessible within the `DELETE`. See Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”.

Single-Table Syntax

```
DELETE [LOW_PRIORITY] [QUICK] [IGNORE] FROM tbl_name
    [PARTITION (partition_name,...)]
    [WHERE where_condition]
    [ORDER BY ...]
    [LIMIT row_count]
```

The `DELETE` statement deletes rows from *`tbl_name`* and returns the number of deleted rows. To check the number of deleted rows, call the `ROW_COUNT()` function described in Section 12.14, “Information Functions”.

Main Clauses

The conditions in the optional `WHERE` clause identify which rows to delete. With no `WHERE` clause, all rows are deleted.

`where_condition` is an expression that evaluates to true for each row to be deleted. It is specified as described in Section 13.2.10, “SELECT Syntax”.

If the `ORDER BY` clause is specified, the rows are deleted in the order that is specified. The `LIMIT` clause places a limit on the number of rows that can be deleted. These clauses apply to single-table deletes, but not multi-table deletes.

Multiple-Table Syntax

```
DELETE [LOW_PRIORITY] [QUICK] [IGNORE]
    tbl_name[.*] [, tbl_name[.*]] ...
    FROM table_references
    [WHERE where_condition]
```

Or:

```
DELETE [LOW_PRIORITY] [QUICK] [IGNORE]
    FROM tbl_name[.*] [, tbl_name[.*]] ...
    USING table_references
    [WHERE where_condition]
```

Privileges

You need the DELETE privilege on a table to delete rows from it. You need only the SELECT privilege for any columns that are only read, such as those named in the WHERE clause.

Performance

When you do not need to know the number of deleted rows, the TRUNCATE TABLE statement is a faster way to empty a table than a DELETE statement with no WHERE clause. Unlike DELETE, TRUNCATE TABLE cannot be used within a transaction or if you have a lock on the table. See Section 13.1.30, “TRUNCATE TABLE Syntax” and Section 13.3.5, “LOCK TABLES and UNLOCK TABLES Syntax”.

The speed of delete operations may also be affected by factors discussed in Section 8.2.5.3, “Optimizing DELETE Statements”.

To ensure that a given DELETE statement does not take too much time, the MySQL-specific LIMIT *row_count* clause for DELETE specifies the maximum number of rows to be deleted. If the number of rows to delete is larger than the limit, repeat the DELETE statement until the number of affected rows is less than the LIMIT value.

Subqueries

You cannot delete from a table and select from the same table in a subquery.

Partitioned Tables

DELETE supports explicit partition selection using the PARTITION option, which takes a comma-separated list of the names of one or more partitions or subpartitions (or both) from which to select rows to be dropped. Partitions not included in the list are ignored. Given a partitioned table t with a partition named p0, executing the statement DELETE FROM t PARTITION (p0) has the same effect on the table as executing ALTER TABLE t TRUNCATE PARTITION (p0); in both cases, all rows in partition p0 are dropped.

PARTITION can be used along with a WHERE condition, in which case the condition is tested only on rows in the listed partitions. For example, DELETE FROM t PARTITION (p0) WHERE c < 5 deletes rows only from partition p0 for which the condition c < 5 is true; rows in any other partitions are not checked and thus not affected by the DELETE.

The PARTITION option can also be used in multiple-table DELETE statements. You can use up to one such option per table named in the FROM option.

See Section 22.5, “Partition Selection”, for more information and examples.

Auto-Increment Columns

If you delete the row containing the maximum value for an AUTO_INCREMENT column, the value is not reused for a MyISAM or InnoDB table. If you delete all rows in the table with DELETE FROM

tbl_name (without a WHERE clause) in autocommit mode, the sequence starts over for all storage engines except InnoDB and MyISAM. There are some exceptions to this behavior for InnoDB tables, as discussed in Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”.

For MyISAM tables, you can specify an AUTO_INCREMENT secondary column in a multiple-column key. In this case, reuse of values deleted from the top of the sequence occurs even for MyISAM tables. See Section 3.6.9, “Using AUTO_INCREMENT”.

Modifiers

The DELETE statement supports the following modifiers:

- If you specify LOW_PRIORITY, the server delays execution of the DELETE until no other clients are reading from the table. This affects only storage engines that use only table-level locking (such as MyISAM, MEMORY, and MERGE).
- For MyISAM tables, if you use the QUICK modifier, the storage engine does not merge index leaves during delete, which may speed up some kinds of delete operations.
- The IGNORE modifier causes MySQL to ignore errors during the process of deleting rows. (Errors encountered during the parsing stage are processed in the usual manner.) Errors that are ignored due to the use of IGNORE are returned as warnings. For more information, see Comparison of the IGNORE Keyword and Strict SQL Mode.

Order of Deletion

If the DELETE statement includes an ORDER BY clause, rows are deleted in the order specified by the clause. This is useful primarily in conjunction with LIMIT. For example, the following statement finds rows matching the WHERE clause, sorts them by timestamp_column, and deletes the first (oldest) one:

```
DELETE FROM somelog WHERE user = 'jcole'
ORDER BY timestamp_column LIMIT 1;
```

ORDER BY also helps to delete rows in an order required to avoid referential integrity violations.

InnoDB Tables

If you are deleting many rows from a large table, you may exceed the lock table size for an InnoDB table. To avoid this problem, or simply to minimize the time that the table remains locked, the following strategy (which does not use DELETE at all) might be helpful:

1. Select the rows *not* to be deleted into an empty table that has the same structure as the original table:

   ```
   INSERT INTO t_copy SELECT * FROM t WHERE ... ;
   ```

2. Use RENAME TABLE to atomically move the original table out of the way and rename the copy to the original name:

   ```
   RENAME TABLE t TO t_old, t_copy TO t;
   ```

3. Drop the original table:

   ```
   DROP TABLE t_old;
   ```

No other sessions can access the tables involved while RENAME TABLE executes, so the rename operation is not subject to concurrency problems. See Section 13.1.29, “RENAME TABLE Syntax”.

MyISAM Tables

In MyISAM tables, deleted rows are maintained in a linked list and subsequent INSERT operations reuse old row positions. To reclaim unused space and reduce file sizes, use the OPTIMIZE TABLE statement or the myisamchk utility to reorganize tables. OPTIMIZE TABLE is easier to use, but myisamchk is faster. See Section 13.7.2.4, “OPTIMIZE TABLE Syntax”, and Section 4.6.4, “myisamchk — MyISAM Table-Maintenance Utility”.

The QUICK modifier affects whether index leaves are merged for delete operations. DELETE QUICK is most useful for applications where index values for deleted rows are replaced by similar index values from rows inserted later. In this case, the holes left by deleted values are reused.

DELETE QUICK is not useful when deleted values lead to underfilled index blocks spanning a range of index values for which new inserts occur again. In this case, use of QUICK can lead to wasted space in the index that remains unreclaimed. Here is an example of such a scenario:

1. Create a table that contains an indexed AUTO_INCREMENT column.
2. Insert many rows into the table. Each insert results in an index value that is added to the high end of the index.
3. Delete a block of rows at the low end of the column range using DELETE QUICK.

In this scenario, the index blocks associated with the deleted index values become underfilled but are not merged with other index blocks due to the use of QUICK. They remain underfilled when new inserts occur, because new rows do not have index values in the deleted range. Furthermore, they remain underfilled even if you later use DELETE without QUICK, unless some of the deleted index values happen to lie in index blocks within or adjacent to the underfilled blocks. To reclaim unused index space under these circumstances, use OPTIMIZE TABLE.

If you are going to delete many rows from a table, it might be faster to use DELETE QUICK followed by OPTIMIZE TABLE. This rebuilds the index rather than performing many index block merge operations.

Multi-Table Deletes

You can specify multiple tables in a DELETE statement to delete rows from one or more tables depending on the condition in the WHERE clause. You cannot use ORDER BY or LIMIT in a multiple-table DELETE. The *table_references* clause lists the tables involved in the join, as described in Section 13.2.10.2, “JOIN Syntax”.

For the first multiple-table syntax, only matching rows from the tables listed before the FROM clause are deleted. For the second multiple-table syntax, only matching rows from the tables listed in the FROM clause (before the USING clause) are deleted. The effect is that you can delete rows from many tables at the same time and have additional tables that are used only for searching:

```
DELETE t1, t2 FROM t1 INNER JOIN t2 INNER JOIN t3
WHERE t1.id=t2.id AND t2.id=t3.id;
```

Or:

```
DELETE FROM t1, t2 USING t1 INNER JOIN t2 INNER JOIN t3
WHERE t1.id=t2.id AND t2.id=t3.id;
```

These statements use all three tables when searching for rows to delete, but delete matching rows only from tables t1 and t2.

The preceding examples use INNER JOIN, but multiple-table DELETE statements can use other types of join permitted in SELECT statements, such as LEFT JOIN. For example, to delete rows that exist in t1 that have no match in t2, use a LEFT JOIN:

```
DELETE t1 FROM t1 LEFT JOIN t2 ON t1.id=t2.id WHERE t2.id IS NULL;
```

The syntax permits `.*` after each *`tbl_name`* for compatibility with `Access`.

If you use a multiple-table `DELETE` statement involving `InnoDB` tables for which there are foreign key constraints, the MySQL optimizer might process tables in an order that differs from that of their parent/child relationship. In this case, the statement fails and rolls back. Instead, you should delete from a single table and rely on the `ON DELETE` capabilities that `InnoDB` provides to cause the other tables to be modified accordingly.

Note

If you declare an alias for a table, you must use the alias when referring to the table:

```
DELETE t1 FROM test AS t1, test2 WHERE ...
```

Table aliases in a multiple-table `DELETE` should be declared only in the *`table_references`* part of the statement. Elsewhere, alias references are permitted but not alias declarations.

Correct:

```
DELETE a1, a2 FROM t1 AS a1 INNER JOIN t2 AS a2
WHERE a1.id=a2.id;

DELETE FROM a1, a2 USING t1 AS a1 INNER JOIN t2 AS a2
WHERE a1.id=a2.id;
```

Incorrect:

```
DELETE t1 AS a1, t2 AS a2 FROM t1 INNER JOIN t2
WHERE a1.id=a2.id;

DELETE FROM t1 AS a1, t2 AS a2 USING t1 INNER JOIN t2
WHERE a1.id=a2.id;
```

13.2.3 DO Syntax

```
DO expr [, expr] ...
```

`DO` executes the expressions but does not return any results. In most respects, `DO` is shorthand for `SELECT expr, ...`, but has the advantage that it is slightly faster when you do not care about the result.

`DO` is useful primarily with functions that have side effects, such as `RELEASE_LOCK()`.

Example: This `SELECT` statement pauses, but also produces a result set:

```
mysql> SELECT SLEEP(5);
+----------+
| SLEEP(5) |
+----------+
|        0 |
+----------+
1 row in set (5.02 sec)
```

`DO`, on the other hand, pauses without producing a result set.:

```
mysql> DO SLEEP(5);
```

```
Query OK, 0 rows affected (4.99 sec)
```

This could be useful, for example in a stored function or trigger, which prohibit statements that produce result sets.

`DO` only executes expressions. It cannot be used in all cases where `SELECT` can be used. For example, `DO id FROM t1` is invalid because it references a table.

13.2.4 HANDLER Syntax

```
HANDLER tbl_name OPEN [ [AS] alias]

HANDLER tbl_name READ index_name { = | <= | >= | < | > } (value1,value2,...)
    [ WHERE where_condition ] [LIMIT ... ]
HANDLER tbl_name READ index_name { FIRST | NEXT | PREV | LAST }
    [ WHERE where_condition ] [LIMIT ... ]
HANDLER tbl_name READ { FIRST | NEXT }
    [ WHERE where_condition ] [LIMIT ... ]

HANDLER tbl_name CLOSE
```

The `HANDLER` statement provides direct access to table storage engine interfaces. It is available for `InnoDB` and `MyISAM` tables.

The `HANDLER ... OPEN` statement opens a table, making it accessible using subsequent `HANDLER ... READ` statements. This table object is not shared by other sessions and is not closed until the session calls `HANDLER ... CLOSE` or the session terminates.

If you open the table using an alias, further references to the open table with other `HANDLER` statements must use the alias rather than the table name. If you do not use an alias, but open the table using a table name qualified by the database name, further references must use the unqualified table name. For example, for a table opened using `mydb.mytable`, further references must use `mytable`.

The first `HANDLER ... READ` syntax fetches a row where the index specified satisfies the given values and the `WHERE` condition is met. If you have a multiple-column index, specify the index column values as a comma-separated list. Either specify values for all the columns in the index, or specify values for a leftmost prefix of the index columns. Suppose that an index `my_idx` includes three columns named `col_a`, `col_b`, and `col_c`, in that order. The `HANDLER` statement can specify values for all three columns in the index, or for the columns in a leftmost prefix. For example:

```
HANDLER ... READ my_idx = (col_a_val,col_b_val,col_c_val) ...
HANDLER ... READ my_idx = (col_a_val,col_b_val) ...
HANDLER ... READ my_idx = (col_a_val) ...
```

To employ the `HANDLER` interface to refer to a table's `PRIMARY KEY`, use the quoted identifier `` `PRIMARY` ``:

```
HANDLER tbl_name READ `PRIMARY` ...
```

The second `HANDLER ... READ` syntax fetches a row from the table in index order that matches the `WHERE` condition.

The third `HANDLER ... READ` syntax fetches a row from the table in natural row order that matches the `WHERE` condition. It is faster than `HANDLER tbl_name READ index_name` when a full table scan is desired. Natural row order is the order in which rows are stored in a `MyISAM` table data file. This statement works for `InnoDB` tables as well, but there is no such concept because there is no separate data file.

Without a `LIMIT` clause, all forms of `HANDLER ... READ` fetch a single row if one is available. To return a specific number of rows, include a `LIMIT` clause. It has the same syntax as for the `SELECT` statement. See Section 13.2.10, “SELECT Syntax”.

`HANDLER ... CLOSE` closes a table that was opened with `HANDLER ... OPEN`.

There are several reasons to use the `HANDLER` interface instead of normal `SELECT` statements:

- `HANDLER` is faster than `SELECT`:
  - A designated storage engine handler object is allocated for the `HANDLER ... OPEN`. The object is reused for subsequent `HANDLER` statements for that table; it need not be reinitialized for each one.
  - There is less parsing involved.
  - There is no optimizer or query-checking overhead.
  - The handler interface does not have to provide a consistent look of the data (for example, dirty reads are permitted), so the storage engine can use optimizations that `SELECT` does not normally permit.
- `HANDLER` makes it easier to port to MySQL applications that use a low-level `ISAM`-like interface. (See Section 15.19, “InnoDB memcached Plugin” for an alternative way to adapt applications that use the key-value store paradigm.)
- `HANDLER` enables you to traverse a database in a manner that is difficult (or even impossible) to accomplish with `SELECT`. The `HANDLER` interface is a more natural way to look at data when working with applications that provide an interactive user interface to the database.

`HANDLER` is a somewhat low-level statement. For example, it does not provide consistency. That is, `HANDLER ... OPEN` does *not* take a snapshot of the table, and does *not* lock the table. This means that after a `HANDLER ... OPEN` statement is issued, table data can be modified (by the current session or other sessions) and these modifications might be only partially visible to `HANDLER ... NEXT` or `HANDLER ... PREV` scans.

An open handler can be closed and marked for reopen, in which case the handler loses its position in the table. This occurs when both of the following circumstances are true:

- Any session executes `FLUSH TABLES` or DDL statements on the handler's table.
- The session in which the handler is open executes non-`HANDLER` statements that use tables.

`TRUNCATE TABLE` for a table closes all handlers for the table that were opened with `HANDLER OPEN`.

If a table is flushed with `FLUSH TABLES tbl_name WITH READ LOCK` was opened with `HANDLER`, the handler is implicitly flushed and loses its position.

13.2.5 IMPORT TABLE Syntax

```
IMPORT TABLE FROM sdi_file [, sdi_file] ...
```

The `IMPORT TABLE` statement imports `MyISAM` tables based on information contained in `.SDI` (Serialized Dictionary Information) metadata files. `IMPORT TABLE` requires the `FILE` privilege to read the `.SDI` and table content files, and the `CREATE` privilege for the table to be created.

Tables can be exported from one server using `mysqldump` to write a file of SQL statements and imported into another server using `mysql` to process the dump file. `IMPORT TABLE` provides a faster alternative using the “raw” table files.

Prior to import, the files that provide the table content must be placed in the appropriate schema directory for the import server, and the `.SDI` file must be located in a directory accessible to the server. For example, the `.SDI` file can be placed in the directory named by the `secure_file_priv` system variable, or (if `secure_file_priv` is empty) in a directory under the server data directory.

The following example describes how to export `MyISAM` tables named `employees` and `managers` from the `hr` schema of one server and import them into the `hr` schema of another server. The example uses these assumptions (to perform a similar operation on your own system, modify the path names as appropriate):

- For the export server, *`export_basedir`* represents its base directory, and its data directory is *`export_basedir`*`/data`.
- For the import server, *`import_basedir`* represents its base directory, and its data directory is *`import_basedir`*`/data`.
- Table files are exported from the export server into the `/tmp/export` directory and this directory is secure (not accessible to other users).
- The import server uses `/tmp/mysql-files` as the directory named by its `secure_file_priv` system variable.

To export tables from the export server, use this procedure:

1. Ensure a consistent snapshot by executing this statement to lock the tables so that they cannot be modified during export:

   ```
   mysql> FLUSH TABLES hr.employees, hr.managers WITH READ LOCK;
   ```

   While the lock is in effect, the tables can still be used, but only for read access.

2. At the file system level, copy the `.SDI` and table content files from the `hr` schema directory to the secure export directory:

   - The `.SDI` file is located in the `hr` schema directory, but might not have exactly the same basename as the table name. For example, the `.SDI` files for the `employees` and `managers` tables might be named `employees_125.SDI` and `managers_238.SDI`.
   - For a `MyISAM` table, the content files are its `.MYD` data file and `.MYI` index file.

   Given those file names, the copy commands look like this:

   ```
   shell> cd export_basedir/data/hr
   shell> cp employees_125.SDI /tmp/export
   shell> cp managers_238.SDI /tmp/export
   shell> cp employees.{MYD,MYI} /tmp/export
   shell> cp managers.{MYD,MYI} /tmp/export
   ```

3. Unlock the tables:

   ```
   mysql> UNLOCK TABLES;
   ```

To import tables into the import server, use this procedure:

1. The import schema must exist. If necessary, execute this statement to create it:

   ```
   mysql> CREATE SCHEMA hr;
   ```

2. At the file system level, copy the `.SDI` files to the import server `secure_file_priv` directory, `/tmp/mysql-files`. Also, copy the table content files to the `hr` schema directory:

   ```
   shell> cd /tmp/export
   shell> cp employees_125.SDI /tmp/mysql-files
   shell> cp managers_238.SDI /tmp/mysql-files
   shell> cp employees.{MYD,MYI} import_basedir/data/hr
   ```

```
shell> cp managers.{MYD,MYI} import_basedir/data/hr
```

3. Import the tables by executing an `IMPORT TABLE` statement that names the `.SDI` files:

```
mysql> IMPORT TABLE FROM
       '/tmp/mysql-files/employees.SDI',
       '/tmp/mysql-files/managers.SDI';
```

The `.SDI` file need not be placed in the import server directory named by the `secure_file_priv` system variable if that variable is empty; it can be in any directory accessible to the server, including the schema directory for the imported table. If the `.SDI` file is placed in that directory, however, it may be rewritten; the import operation creates a new `.SDI` file for the table, which will overwrite the old `.SDI` file if the operation uses the same file name for the new file.

Each *`sdi_file`* value must be a string literal that names the `.SDI` file for a table or is a pattern that matches `.SDI` files. If the string is a pattern, any leading directory path and the `.SDI` file name suffix must be given literally. Pattern characters are permitted only in the base name part of the file name:

- `?` matches any single character
- `*` matches any sequence of characters, including no characters

Using a pattern, the previous `IMPORT TABLE` statement could have been written like this (assuming that the `/tmp/mysql-files` directory contains no other `.SDI` files matching the pattern):

```
IMPORT TABLE FROM '/tmp/mysql-files/*.SDI';
```

To interpret the location of `.SDI` file path names, the server uses the same rules for `IMPORT TABLE` as the server-side rules for `LOAD DATA` (that is, the non-`LOCAL` rules). See Section 13.2.7, “LOAD DATA INFILE Syntax”, paying particular attention to the rules used to interpret relative path names.

`IMPORT TABLE` fails if the `.SDI` or table files cannot be located. After importing a table, the server attempts to open it and reports as warnings any problems detected. To attempt a repair to correct any reported issues, use `REPAIR TABLE`.

`IMPORT TABLE` is not written to the binary log.

Restrictions and Limitations

`IMPORT TABLE` applies only to non-`TEMPORARY` `MyISAM` tables. It does not apply to tables created with a transactional storage engine, tables created with `CREATE TEMPORARY TABLE`, or views.

The table data and index files must be placed in the schema directory for the import server prior to the import operation, unless the table as defined on the export server uses the `DATA DIRECTORY` or `INDEX DIRECTORY` table options. In that case, modify the import procedure using one of these alternatives before executing the `IMPORT TABLE` statement:

- Put the data and index files into the same directory on the import server host as on the export server host, and create symlinks in the import server schema directory to those files.
- Put the data and index files into an import server host directory different from that on the export server host, and create symlinks in the import server schema directory to those files. In addition, modify the `.SDI` file to reflect the different file locations.
- Put the data and index files into the schema directory on the import server host, and modify the `.SDI` file to remove the data and index directory table options.

Any collation IDs stored in the `.SDI` file must refer to the same collations on the export and import servers.

Trigger information for a table is not serialized into the table `.SDI` file, so triggers are not restored by the import operation.

Some edits to an `.SDI` file are permissible prior to executing the `IMPORT TABLE` statement, whereas others are problematic or may even cause the import operation to fail:

- Changing the data directory and index directory table options is required if the locations of the data and index files differ between the export and import servers.
- Changing the schema name is required to import the table into a different schema on the import server than on the export server.
- Changing schema and table names may be required to accommodate differences between file system case-sensitivity semantics on the export and import servers or differences in `lower_case_table_names` settings. Changing the table names in the `.SDI` file may require renamimg the table files as well.
- In some cases, changes to column definitions are permitted. Changing data types is likely to cause problems.

13.2.6 INSERT Syntax

```
INSERT [LOW_PRIORITY | DELAYED | HIGH_PRIORITY] [IGNORE]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    [(col_name,...)]
    {VALUES | VALUE} ({expr | DEFAULT},...),(...),...
    [ ON DUPLICATE KEY UPDATE
      col_name=expr
        [, col_name=expr] ... ]
```

Or:

```
INSERT [LOW_PRIORITY | DELAYED | HIGH_PRIORITY] [IGNORE]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    SET col_name={expr | DEFAULT}, ...
    [ ON DUPLICATE KEY UPDATE
      col_name=expr
        [, col_name=expr] ... ]
```

Or:

```
INSERT [LOW_PRIORITY | HIGH_PRIORITY] [IGNORE]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    [(col_name,...)]
    SELECT ...
    [ ON DUPLICATE KEY UPDATE
      col_name=expr
        [, col_name=expr] ... ]
```

`INSERT` inserts new rows into an existing table. The `INSERT ... VALUES` and `INSERT ... SET` forms of the statement insert rows based on explicitly specified values. The `INSERT ... SELECT` form inserts rows selected from another table or tables. `INSERT ... SELECT` is discussed further in Section 13.2.6.1, “INSERT ... SELECT Syntax”.

When inserting into a partitioned table, you can control which partitions and subpartitions accept new rows. The `PARTITION` option takes a comma-separated list of the names of one or more partitions or subpartitions (or both) of the table. If any of the rows to be inserted by a given `INSERT` statement do not match one of the partitions listed, the `INSERT` statement fails with the error `Found a row not matching the given partition set`. See Section 22.5, “Partition Selection”, for more information and examples.

In MySQL 8.0, the DELAYED keyword is accepted but ignored by the server. See Section 13.2.6.2, “INSERT DELAYED Syntax”, for the reasons for this.

You can use REPLACE instead of INSERT to overwrite old rows. REPLACE is the counterpart to INSERT IGNORE in the treatment of new rows that contain unique key values that duplicate old rows: The new rows are used to replace the old rows rather than being discarded. See Section 13.2.9, “REPLACE Syntax”.

tbl_name is the table into which rows should be inserted. The columns for which the statement provides values can be specified as follows:

- You can provide a comma-separated list of column names following the table name. In this case, a value for each named column must be provided by the VALUES list or the SELECT statement.
- If you do not specify a list of column names for INSERT ... VALUES or INSERT ... SELECT, values for every column in the table must be provided by the VALUES list or the SELECT statement. If you do not know the order of the columns in the table, use DESCRIBE *tbl_name* to find out.
- The SET clause indicates the column names explicitly.

Column values can be given in several ways:

- If you are not running in strict SQL mode, any column not explicitly given a value is set to its default (explicit or implicit) value. For example, if you specify a column list that does not name all the columns in the table, unnamed columns are set to their default values. Default value assignment is described in Section 11.7, “Data Type Default Values”. See also Section 1.8.3.3, “Constraints on Invalid Data”.

  If you want an INSERT statement to generate an error unless you explicitly specify values for all columns that do not have a default value, you should use strict mode. See Section 5.1.8, “Server SQL Modes”.

- Use the keyword DEFAULT to set a column explicitly to its default value. This makes it easier to write INSERT statements that assign values to all but a few columns, because it enables you to avoid writing an incomplete VALUES list that does not include a value for each column in the table. Otherwise, you would have to write out the list of column names corresponding to each value in the VALUES list.

  You can also use DEFAULT(*col_name*) as a more general form that can be used in expressions to produce a given column's default value.

- If both the column list and the VALUES list are empty, INSERT creates a row with each column set to its default value:

```
INSERT INTO tbl_name () VALUES();
```

  In strict mode, an error occurs if any column doesn't have a default value. Otherwise, MySQL uses the implicit default value for any column that does not have an explicitly defined default.

- You can specify an expression *expr* to provide a column value. This might involve type conversion if the type of the expression does not match the type of the column, and conversion of a given value can result in different inserted values depending on the data type. For example, inserting the string '1999.0e-2' into an INT, FLOAT, DECIMAL(10,6), or YEAR column results in the values 1999, 19.9921, 19.992100, and 1999 being inserted, respectively. The reason the value stored in the INT and YEAR columns is 1999 is that the string-to-integer conversion looks only at as much of the initial part of the string as may be considered a valid integer or year. For the floating-point and fixed-point columns, the string-to-floating-point conversion considers the entire string a valid floating-point value.

  An expression *expr* can refer to any column that was set earlier in a value list. For example, you can do this because the value for col2 refers to col1, which has previously been assigned:

```
INSERT INTO tbl_name (col1,col2) VALUES(15,col1*2);
```

But the following is not legal, because the value for `col1` refers to `col2`, which is assigned after `col1`:

```
INSERT INTO tbl_name (col1,col2) VALUES(col2*2,15);
```

One exception involves columns that contain `AUTO_INCREMENT` values. Because the `AUTO_INCREMENT` value is generated after other value assignments, any reference to an `AUTO_INCREMENT` column in the assignment returns a `0`.

`INSERT` statements that use `VALUES` syntax can insert multiple rows. To do this, include multiple lists of column values, each enclosed within parentheses and separated by commas. Example:

```
INSERT INTO tbl_name (a,b,c) VALUES(1,2,3),(4,5,6),(7,8,9);
```

The values list for each row must be enclosed within parentheses. The following statement is illegal because the number of values in the list does not match the number of column names:

```
INSERT INTO tbl_name (a,b,c) VALUES(1,2,3,4,5,6,7,8,9);
```

`VALUE` is a synonym for `VALUES` in this context. Neither implies anything about the number of values lists, and either may be used whether there is a single values list or multiple lists.

The affected-rows value for an `INSERT` can be obtained using the `ROW_COUNT()` function (see Section 12.14, “Information Functions”), or the `mysql_affected_rows()` C API function (see Section 27.7.7.1, “mysql_affected_rows()”).

If you use an `INSERT ... VALUES` statement with multiple value lists or `INSERT ... SELECT`, the statement returns an information string in this format:

```
Records: 100 Duplicates: 0 Warnings: 0
```

`Records` indicates the number of rows processed by the statement. (This is not necessarily the number of rows actually inserted because `Duplicates` can be nonzero.) `Duplicates` indicates the number of rows that could not be inserted because they would duplicate some existing unique index value. `Warnings` indicates the number of attempts to insert column values that were problematic in some way. Warnings can occur under any of the following conditions:

- Inserting `NULL` into a column that has been declared `NOT NULL`. For multiple-row `INSERT` statements or `INSERT INTO ... SELECT` statements, the column is set to the implicit default value for the column data type. This is `0` for numeric types, the empty string (`''`) for string types, and the “zero” value for date and time types. `INSERT INTO ... SELECT` statements are handled the same way as multiple-row inserts because the server does not examine the result set from the `SELECT` to see whether it returns a single row. (For a single-row `INSERT`, no warning occurs when `NULL` is inserted into a `NOT NULL` column. Instead, the statement fails with an error.)

- Setting a numeric column to a value that lies outside the column's range. The value is clipped to the closest endpoint of the range.

- Assigning a value such as `'10.34 a'` to a numeric column. The trailing nonnumeric text is stripped off and the remaining numeric part is inserted. If the string value has no leading numeric part, the column is set to `0`.

- Inserting a string into a string column (`CHAR`, `VARCHAR`, `TEXT`, or `BLOB`) that exceeds the column's maximum length. The value is truncated to the column's maximum length.

- Inserting a value into a date or time column that is illegal for the data type. The column is set to the appropriate zero value for the type.

If a generated column is inserted into explicitly, the only permitted value is `DEFAULT`. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

If you are using the C API, the information string can be obtained by invoking the `mysql_info()` function. See Section 27.7.7.36, “mysql_info()”.

If `INSERT` inserts a row into a table that has an `AUTO_INCREMENT` column, you can find the value used for that column by using the SQL `LAST_INSERT_ID()` function. From within the C API, use the `mysql_insert_id()` function. However, you should note that the two functions do not always behave identically. The behavior of `INSERT` statements with respect to `AUTO_INCREMENT` columns is discussed further in Section 12.14, “Information Functions”, and Section 27.7.7.38, “mysql_insert_id()”.

The `INSERT` statement supports the following modifiers:

- `INSERT DELAYED` was deprecated in MySQL 5.7, and is scheduled for eventual removal. In MySQL 8.0, the `DELAYED` modifier is accepted but ignored. Use `INSERT` (without `DELAYED`) instead. See Section 13.2.6.2, “INSERT DELAYED Syntax”.

- If you use the `LOW_PRIORITY` modifier, execution of the `INSERT` is delayed until no other clients are reading from the table. This includes other clients that began reading while existing clients are reading, and while the `INSERT LOW_PRIORITY` statement is waiting. It is possible, therefore, for a client that issues an `INSERT LOW_PRIORITY` statement to wait for a very long time.

  Note

  `LOW_PRIORITY` should normally not be used with `MyISAM` tables because doing so disables concurrent inserts. See Section 8.11.3, “Concurrent Inserts”.

  If you specify `HIGH_PRIORITY`, it overrides the effect of the `--low-priority-updates` option if the server was started with that option. It also causes concurrent inserts not to be used. See Section 8.11.3, “Concurrent Inserts”.

  `LOW_PRIORITY` and `HIGH_PRIORITY` affect only storage engines that use only table-level locking (such as `MyISAM`, `MEMORY`, and `MERGE`).

- If you use the `IGNORE` modifier, errors that occur while executing the `INSERT` statement are ignored. For example, without `IGNORE`, a row that duplicates an existing `UNIQUE` index or `PRIMARY KEY` value in the table causes a duplicate-key error and the statement is aborted. With `IGNORE`, the row is discarded and no error occurs. Ignored errors generate warnings instead.

  `IGNORE` has a similar effect on inserts into partitioned tables where no partition matching a given value is found. Without `IGNORE`, such `INSERT` statements are aborted with an error; however, when `INSERT IGNORE` is used, the insert operation fails silently for the row containing the unmatched value, but any rows that are matched are inserted. For an example, see Section 22.2.2, “LIST Partitioning”.

  Data conversions that would trigger errors abort the statement if `IGNORE` is not specified. With `IGNORE`, invalid values are adjusted to the closest values and inserted; warnings are produced but the statement does not abort. You can determine with the `mysql_info()` C API function how many rows were actually inserted into the table.

  For more information, see Comparison of the IGNORE Keyword and Strict SQL Mode.

- If you specify `ON DUPLICATE KEY UPDATE`, and a row is inserted that would cause a duplicate value in a `UNIQUE` index or `PRIMARY KEY`, an `UPDATE` of the old row is performed. The affected-rows value per row is 1 if the row is inserted as a new row, 2 if an existing row is updated, and

0 if an existing row is set to its current values. If you specify the `CLIENT_FOUND_ROWS` flag to `mysql_real_connect()` when connecting to `mysqld`, the affected-rows value is 1 (not 0) if an existing row is set to its current values. See Section 13.2.6.3, “INSERT ... ON DUPLICATE KEY UPDATE Syntax”.

Inserting into a table requires the `INSERT` privilege for the table. If the `ON DUPLICATE KEY UPDATE` clause is used and a duplicate key causes an `UPDATE` to be performed instead, the statement requires the `UPDATE` privilege for the columns to be updated. For columns that are read but not modified you need only the `SELECT` privilege (such as for a column referenced only on the right hand side of an *`col_name`*=*`expr`* assignment in an `ON DUPLICATE KEY UPDATE` clause).

In MySQL 8.0, an `INSERT` statement affecting a partitioned table using a storage engine such as `MyISAM` that employs table-level locks locks only those partitions into which rows are actually inserted. (For storage engines such as `InnoDB` that employ row-level locking, no locking of partitions takes place.) For more information, see Partitioning and Locking.

13.2.6.1 INSERT ... SELECT Syntax

```
INSERT [LOW_PRIORITY | HIGH_PRIORITY] [IGNORE]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    [(col_name,...)]
    SELECT ...
    [ ON DUPLICATE KEY UPDATE col_name=expr, ... ]
```

With `INSERT ... SELECT`, you can quickly insert many rows into a table from one or many tables. For example:

```
INSERT INTO tbl_temp2 (fld_id)
  SELECT tbl_temp1.fld_order_id
  FROM tbl_temp1 WHERE tbl_temp1.fld_order_id > 100;
```

The following conditions hold for a `INSERT ... SELECT` statements:

- Specify `IGNORE` to ignore rows that would cause duplicate-key violations.
- The target table of the `INSERT` statement may appear in the `FROM` clause of the `SELECT` part of the query. (This was not possible in some older versions of MySQL.) However, you cannot insert into a table and select from the same table in a subquery.

  When selecting from and inserting into a table at the same time, MySQL creates a temporary table to hold the rows from the `SELECT` and then inserts those rows into the target table. However, it remains true that you cannot use `INSERT INTO t ... SELECT ... FROM t` when `t` is a `TEMPORARY` table, because `TEMPORARY` tables cannot be referred to twice in the same statement (see Section B.5.6.2, “TEMPORARY Table Problems”).
- `AUTO_INCREMENT` columns work as usual.
- To ensure that the binary log can be used to re-create the original tables, MySQL does not permit concurrent inserts for `INSERT ... SELECT` statements (see Section 8.11.3, “Concurrent Inserts”).
- To avoid ambiguous column reference problems when the `SELECT` and the `INSERT` refer to the same table, provide a unique alias for each table used in the `SELECT` part, and qualify column names in that part with the appropriate alias.

You can explicitly select which partitions or subpartitions (or both) of the source or target table (or both) are to be used with a `PARTITION` option following the name of the table. When `PARTITION` is used with the name of the source table in the `SELECT` portion of the statement, rows are selected only from the partitions or subpartitions named in its partition list. When `PARTITION` is used with the name of the target table for the `INSERT` portion of the statement, then it must be possible to insert all rows selected

into the partitions or subpartitions named in the partition list following the option, else the `INSERT ... SELECT` statement fails. For more information and examples, see Section 22.5, “Partition Selection”.

In the values part of `ON DUPLICATE KEY UPDATE`, you can refer to columns in other tables, as long as you do not use `GROUP BY` in the `SELECT` part. One side effect is that you must qualify nonunique column names in the values part.

The order in which rows are returned by a `SELECT` statement with no `ORDER BY` clause is not determined. This means that, when using replication, there is no guarantee that such a `SELECT` returns rows in the same order on the master and the slave; this can lead to inconsistencies between them. To prevent this from occurring, you should always write `INSERT ... SELECT` statements that are to be replicated as `INSERT ... SELECT ... ORDER BY column`. The choice of `column` does not matter as long as the same order for returning the rows is enforced on both the master and the slave. See also Section 18.4.1.17, “Replication and LIMIT”.

Due to this issue, `INSERT ... SELECT ON DUPLICATE KEY UPDATE` and `INSERT IGNORE ... SELECT` statements are flagged as unsafe for statement-based replication. With this change, such statements produce a warning in the log when using statement-based mode and are logged using the row-based format when using `MIXED` mode. (Bug #11758262, Bug #50439)

See also Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

In MySQL 8.0, an `INSERT ... SELECT` statement that acted on partitioned tables using a storage engine such as `MyISAM` that employs table-level locks locks all partitions of the target table; however, only those partitions that are actually read from the source table are locked. (This does not occur with tables using storage engines such as `InnoDB` that employ row-level locking.) See Partitioning and Locking, for more information.

13.2.6.2 INSERT DELAYED Syntax

```
INSERT DELAYED ...
```

The `DELAYED` option for the `INSERT` statement is a MySQL extension to standard SQL. In previous versions of MySQL, it can be used for certain kinds of tables (such as `MyISAM`), such that when a client uses `INSERT DELAYED`, it gets an okay from the server at once, and the row is queued to be inserted when the table is not in use by any other thread.

`DELAYED` inserts and replaces were deprecated in MySQL 5.6.6. In MySQL 8.0, `DELAYED` is not supported. The server recognizes but ignores the `DELAYED` keyword, handles the insert as a nondelayed insert, and generates an `ER_WARN_LEGACY_SYNTAX_CONVERTED` warning (“INSERT DELAYED is no longer supported. The statement was converted to INSERT”). The `DELAYED` keyword is scheduled for removal in a future release.

13.2.6.3 INSERT ... ON DUPLICATE KEY UPDATE Syntax

If you specify `ON DUPLICATE KEY UPDATE`, and a row is inserted that would cause a duplicate value in a `UNIQUE` index or `PRIMARY KEY`, MySQL performs an `UPDATE` of the old row. For example, if column `a` is declared as `UNIQUE` and contains the value `1`, the following two statements have similar effect:

```
INSERT INTO table (a,b,c) VALUES (1,2,3)
  ON DUPLICATE KEY UPDATE c=c+1;

UPDATE table SET c=c+1 WHERE a=1;
```

(The effects are not identical for an `InnoDB` table where `a` is an auto-increment column. With an auto-increment column, an `INSERT` statement increases the auto-increment value but `UPDATE` does not.)

The `ON DUPLICATE KEY UPDATE` clause can contain multiple column assignments, separated by commas.

With `ON DUPLICATE KEY UPDATE`, the affected-rows value per row is 1 if the row is inserted as a new row, 2 if an existing row is updated, and 0 if an existing row is set to its current values. If you specify the `CLIENT_FOUND_ROWS` flag to `mysql_real_connect()` when connecting to `mysqld`, the affected-rows value is 1 (not 0) if an existing row is set to its current values.

If column `b` is also unique, the `INSERT` is equivalent to this `UPDATE` statement instead:

```
UPDATE table SET c=c+1 WHERE a=1 OR b=2 LIMIT 1;
```

If `a=1 OR b=2` matches several rows, only *one* row is updated. In general, you should try to avoid using an `ON DUPLICATE KEY UPDATE` clause on tables with multiple unique indexes.

You can use the `VALUES(col_name)` function in the `UPDATE` clause to refer to column values from the `INSERT` portion of the `INSERT ... ON DUPLICATE KEY UPDATE` statement. In other words, `VALUES(col_name)` in the `ON DUPLICATE KEY UPDATE` clause refers to the value of `col_name` that would be inserted, had no duplicate-key conflict occurred. This function is especially useful in multiple-row inserts. The `VALUES()` function is meaningful only in `INSERT ... UPDATE` statements and returns `NULL` otherwise. Example:

```
INSERT INTO table (a,b,c) VALUES (1,2,3),(4,5,6)
  ON DUPLICATE KEY UPDATE c=VALUES(a)+VALUES(b);
```

That statement is identical to the following two statements:

```
INSERT INTO table (a,b,c) VALUES (1,2,3)
  ON DUPLICATE KEY UPDATE c=3;
INSERT INTO table (a,b,c) VALUES (4,5,6)
  ON DUPLICATE KEY UPDATE c=9;
```

If a table contains an `AUTO_INCREMENT` column and `INSERT ... ON DUPLICATE KEY UPDATE` inserts or updates a row, the `LAST_INSERT_ID()` function returns the `AUTO_INCREMENT` value.

The `DELAYED` option is ignored when you use `ON DUPLICATE KEY UPDATE`.

Because the results of `INSERT ... SELECT` statements depend on the ordering of rows from the `SELECT` and this order cannot always be guaranteed, it is possible when logging `INSERT ... SELECT ON DUPLICATE KEY UPDATE` statements for the master and the slave to diverge. Thus, `INSERT ... SELECT ON DUPLICATE KEY UPDATE` statements are flagged as unsafe for statement-based replication. With this change, such statements produce a warning in the log when using statement-based mode and are logged using the row-based format when using `MIXED` mode. In addition, an `INSERT ... ON DUPLICATE KEY UPDATE` statement against a table having more than one unique or primary key is also marked as unsafe. (Bug #11765650, Bug #58637) See also Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

In MySQL 8.0, an `INSERT ... ON DUPLICATE KEY UPDATE` on a partitioned table using a storage engine such as `MyISAM` that employs table-level locks locks any partitions of the table in which a partitioning key column is updated. (This does not occur with tables using storage engines such as `InnoDB` that employ row-level locking.) See Partitioning and Locking, for more information.

13.2.7 LOAD DATA INFILE Syntax

```
LOAD DATA [LOW_PRIORITY | CONCURRENT] [LOCAL] INFILE 'file_name'
    [REPLACE | IGNORE]
    INTO TABLE tbl_name
```

```
    [PARTITION (partition_name,...)]
    [CHARACTER SET charset_name]
    [{FIELDS | COLUMNS}
        [TERMINATED BY 'string']
        [[OPTIONALLY] ENCLOSED BY 'char']
        [ESCAPED BY 'char']
    ]
    [LINES
        [STARTING BY 'string']
        [TERMINATED BY 'string']
    ]
    [IGNORE number {LINES | ROWS}]
    [(col_name_or_user_var,...)]
    [SET col_name = expr,...]
```

The `LOAD DATA INFILE` statement reads rows from a text file into a table at a very high speed. `LOAD DATA INFILE` is the complement of `SELECT ... INTO OUTFILE`. (See Section 13.2.10.1, “SELECT ... INTO Syntax”.) To write data from a table to a file, use `SELECT ... INTO OUTFILE`. To read the file back into a table, use `LOAD DATA INFILE`. The syntax of the `FIELDS` and `LINES` clauses is the same for both statements. Both clauses are optional, but `FIELDS` must precede `LINES` if both are specified.

You can also load data files by using the `mysqlimport` utility; it operates by sending a `LOAD DATA INFILE` statement to the server. The `--local` option causes `mysqlimport` to read data files from the client host. You can specify the `--compress` option to get better performance over slow networks if the client and server support the compressed protocol. See Section 4.5.5, “`mysqlimport` — A Data Import Program”.

For more information about the efficiency of `INSERT` versus `LOAD DATA INFILE` and speeding up `LOAD DATA INFILE`, see Section 8.2.5.1, “Optimizing INSERT Statements”.

The file name must be given as a literal string. On Windows, specify backslashes in path names as forward slashes or doubled backslashes. The `character_set_filesystem` system variable controls the interpretation of the file name.

`LOAD DATA` supports explicit partition selection using the `PARTITION` option with a comma-separated list of one or more names of partitions, subpartitions, or both. When this option is used, if any rows from the file cannot be inserted into any of the partitions or subpartitions named in the list, the statement fails with the error `Found a row not matching the given partition set`. For more information, see Section 22.5, “Partition Selection”.

For partitioned tables using storage engines that employ table locks, such as `MyISAM`, `LOAD DATA` cannot prune any partition locks. This does not apply to tables using storage engines which employ row-level locking, such as `InnoDB`. For more information, see Partitioning and Locking.

The server uses the character set indicated by the `character_set_database` system variable to interpret the information in the file. `SET NAMES` and the setting of `character_set_client` do not affect interpretation of input. If the contents of the input file use a character set that differs from the default, it is usually preferable to specify the character set of the file by using the `CHARACTER SET` clause. A character set of `binary` specifies “no conversion.”

`LOAD DATA INFILE` interprets all fields in the file as having the same character set, regardless of the data types of the columns into which field values are loaded. For proper interpretation of file contents, you must ensure that it was written with the correct character set. For example, if you write a data file with `mysqldump -T` or by issuing a `SELECT ... INTO OUTFILE` statement in `mysql`, be sure to use a `--default-character-set` option so that output is written in the character set to be used when the file is loaded with `LOAD DATA INFILE`.

Note

It is not possible to load data files that use the `ucs2`, `utf16`, `utf16le`, or `utf32` character set.

If you use LOW_PRIORITY, execution of the LOAD DATA statement is delayed until no other clients are reading from the table. This affects only storage engines that use only table-level locking (such as MyISAM, MEMORY, and MERGE).

If you specify CONCURRENT with a MyISAM table that satisfies the condition for concurrent inserts (that is, it contains no free blocks in the middle), other threads can retrieve data from the table while LOAD DATA is executing. This option affects the performance of LOAD DATA a bit, even if no other thread is using the table at the same time.

With row-based replication, CONCURRENT is replicated regardless of MySQL version. With statement-based replication CONCURRENT is not replicated prior to MySQL 5.5.1 (see Bug #34628). For more information, see Section 18.4.1.18, “Replication and LOAD DATA INFILE”.

The LOCAL keyword affects expected location of the file and error handling, as described later. LOCAL works only if your server and your client both have been configured to permit it. For example, if mysqld was started with the local_infile system variable disabled, LOCAL does not work. See Section 6.1.6, “Security Issues with LOAD DATA LOCAL”.

The LOCAL keyword affects where the file is expected to be found:

- If LOCAL is specified, the file is read by the client program on the client host and sent to the server. The file can be given as a full path name to specify its exact location. If given as a relative path name, the name is interpreted relative to the directory in which the client program was started.

  When using LOCAL with LOAD DATA, a copy of the file is created in the server's temporary directory. This is *not* the directory determined by the value of tmpdir or slave_load_tmpdir, but rather the operating system's temporary directory, and is not configurable in the MySQL Server. (Typically the system temporary directory is /tmp on Linux systems and C:\WINDOWS\TEMP on Windows.) Lack of sufficient space for the copy in this directory can cause the LOAD DATA LOCAL statement to fail.

- If LOCAL is not specified, the file must be located on the server host and is read directly by the server. The server uses the following rules to locate the file:

  - If the file name is an absolute path name, the server uses it as given.
  - If the file name is a relative path name with one or more leading components, the server searches for the file relative to the server's data directory.
  - If a file name with no leading components is given, the server looks for the file in the database directory of the default database.

In the non-LOCAL case, these rules mean that a file named as ./myfile.txt is read from the server's data directory, whereas the file named as myfile.txt is read from the database directory of the default database. For example, if db1 is the default database, the following LOAD DATA statement reads the file data.txt from the database directory for db1, even though the statement explicitly loads the file into a table in the db2 database:

```
LOAD DATA INFILE 'data.txt' INTO TABLE db2.my_table;
```

Note

The server also use the non-LOCAL rules to locate .SDI files for the IMPORT TABLE statement.

Non-LOCAL load operations read text files located on the server. For security reasons, such operations require that you have the FILE privilege. See Section 6.2.1, “Privileges Provided by MySQL”. Also, non-LOCAL load operations are subject to the secure_file_priv system variable setting. If the variable value is a nonempty directory name, the file to be loaded must be located in that directory. If the variable value is empty (which is insecure), the file need only be readable by the server.

Using LOCAL is a bit slower than letting the server access the files directly, because the contents of the file must be sent over the connection by the client to the server. On the other hand, you do not need the FILE privilege to load local files.

LOCAL also affects error handling:

- With LOAD DATA INFILE, data-interpretation and duplicate-key errors terminate the operation.
- With LOAD DATA LOCAL INFILE, data-interpretation and duplicate-key errors become warnings and the operation continues because the server has no way to stop transmission of the file in the middle of the operation. For duplicate-key errors, this is the same as if IGNORE is specified. IGNORE is explained further later in this section.

The REPLACE and IGNORE keywords control handling of input rows that duplicate existing rows on unique key values:

- If you specify REPLACE, input rows replace existing rows. In other words, rows that have the same value for a primary key or unique index as an existing row. See Section 13.2.9, “REPLACE Syntax”.
- If you specify IGNORE, rows that duplicate an existing row on a unique key value are discarded. For more information, see Comparison of the IGNORE Keyword and Strict SQL Mode.
- If you do not specify either option, the behavior depends on whether the LOCAL keyword is specified. Without LOCAL, an error occurs when a duplicate key value is found, and the rest of the text file is ignored. With LOCAL, the default behavior is the same as if IGNORE is specified; this is because the server has no way to stop transmission of the file in the middle of the operation.

To ignore foreign key constraints during the load operation, issue a SET foreign_key_checks = 0 statement before executing LOAD DATA.

If you use LOAD DATA INFILE on an empty MyISAM table, all nonunique indexes are created in a separate batch (as for REPAIR TABLE). Normally, this makes LOAD DATA INFILE much faster when you have many indexes. In some extreme cases, you can create the indexes even faster by turning them off with ALTER TABLE ... DISABLE KEYS before loading the file into the table and using ALTER TABLE ... ENABLE KEYS to re-create the indexes after loading the file. See Section 8.2.5.1, “Optimizing INSERT Statements”.

For both the LOAD DATA INFILE and SELECT ... INTO OUTFILE statements, the syntax of the FIELDS and LINES clauses is the same. Both clauses are optional, but FIELDS must precede LINES if both are specified.

If you specify a FIELDS clause, each of its subclauses (TERMINATED BY, [OPTIONALLY] ENCLOSED BY, and ESCAPED BY) is also optional, except that you must specify at least one of them.

If you specify no FIELDS or LINES clause, the defaults are the same as if you had written this:

```
FIELDS TERMINATED BY '\t' ENCLOSED BY '' ESCAPED BY '\\'
LINES TERMINATED BY '\n' STARTING BY ''
```

(Backslash is the MySQL escape character within strings in SQL statements, so to specify a literal backslash, you must specify two backslashes for the value to be interpreted as a single backslash. The escape sequences '\t' and '\n' specify tab and newline characters, respectively.)

In other words, the defaults cause LOAD DATA INFILE to act as follows when reading input:

- Look for line boundaries at newlines.
- Do not skip over any line prefix.
- Break lines into fields at tabs.

- Do not expect fields to be enclosed within any quoting characters.
- Interpret characters preceded by the escape character `\` as escape sequences. For example, `\t`, `\n`, and `\\` signify tab, newline, and backslash, respectively. See the discussion of `FIELDS ESCAPED BY` later for the full list of escape sequences.

Conversely, the defaults cause `SELECT ... INTO OUTFILE` to act as follows when writing output:

- Write tabs between fields.
- Do not enclose fields within any quoting characters.
- Use `\` to escape instances of tab, newline, or `\` that occur within field values.
- Write newlines at the ends of lines.

Note

If you have generated the text file on a Windows system, you might have to use `LINES TERMINATED BY '\r\n'` to read the file properly, because Windows programs typically use two characters as a line terminator. Some programs, such as `WordPad`, might use `\r` as a line terminator when writing files. To read such files, use `LINES TERMINATED BY '\r'`.

If all the lines you want to read in have a common prefix that you want to ignore, you can use `LINES STARTING BY 'prefix_string'` to skip over the prefix, *and anything before it*. If a line does not include the prefix, the entire line is skipped. Suppose that you issue the following statement:

```
LOAD DATA INFILE '/tmp/test.txt' INTO TABLE test
  FIELDS TERMINATED BY ','  LINES STARTING BY 'xxx';
```

If the data file looks like this:

```
xxx"abc",1
something xxx"def",2
"ghi",3
```

The resulting rows will be `("abc",1)` and `("def",2)`. The third row in the file is skipped because it does not contain the prefix.

The `IGNORE number LINES` option can be used to ignore lines at the start of the file. For example, you can use `IGNORE 1 LINES` to skip over an initial header line containing column names:

```
LOAD DATA INFILE '/tmp/test.txt' INTO TABLE test IGNORE 1 LINES;
```

When you use `SELECT ... INTO OUTFILE` in tandem with `LOAD DATA INFILE` to write data from a database into a file and then read the file back into the database later, the field- and line-handling options for both statements must match. Otherwise, `LOAD DATA INFILE` will not interpret the contents of the file properly. Suppose that you use `SELECT ... INTO OUTFILE` to write a file with fields delimited by commas:

```
SELECT * INTO OUTFILE 'data.txt'
  FIELDS TERMINATED BY ','
  FROM table2;
```

To read the comma-delimited file back in, the correct statement would be:

```
LOAD DATA INFILE 'data.txt' INTO TABLE table2
```

```
  FIELDS TERMINATED BY ',';
```

If instead you tried to read in the file with the statement shown following, it wouldn't work because it instructs `LOAD DATA INFILE` to look for tabs between fields:

```
LOAD DATA INFILE 'data.txt' INTO TABLE table2
  FIELDS TERMINATED BY '\t';
```

The likely result is that each input line would be interpreted as a single field.

`LOAD DATA INFILE` can be used to read files obtained from external sources. For example, many programs can export data in comma-separated values (CSV) format, such that lines have fields separated by commas and enclosed within double quotation marks, with an initial line of column names. If the lines in such a file are terminated by carriage return/newline pairs, the statement shown here illustrates the field- and line-handling options you would use to load the file:

```
LOAD DATA INFILE 'data.txt' INTO TABLE tbl_name
  FIELDS TERMINATED BY ',' ENCLOSED BY '"'
  LINES TERMINATED BY '\r\n'
  IGNORE 1 LINES;
```

If the input values are not necessarily enclosed within quotation marks, use `OPTIONALLY` before the `ENCLOSED BY` keywords.

Any of the field- or line-handling options can specify an empty string (`''`). If not empty, the `FIELDS [OPTIONALLY] ENCLOSED BY` and `FIELDS ESCAPED BY` values must be a single character. The `FIELDS TERMINATED BY`, `LINES STARTING BY`, and `LINES TERMINATED BY` values can be more than one character. For example, to write lines that are terminated by carriage return/linefeed pairs, or to read a file containing such lines, specify a `LINES TERMINATED BY '\r\n'` clause.

To read a file containing jokes that are separated by lines consisting of `%%`, you can do this

```
CREATE TABLE jokes
  (a INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
  joke TEXT NOT NULL);
LOAD DATA INFILE '/tmp/jokes.txt' INTO TABLE jokes
  FIELDS TERMINATED BY ''
  LINES TERMINATED BY '\n%%\n' (joke);
```

`FIELDS [OPTIONALLY] ENCLOSED BY` controls quoting of fields. For output (`SELECT ... INTO OUTFILE`), if you omit the word `OPTIONALLY`, all fields are enclosed by the `ENCLOSED BY` character. An example of such output (using a comma as the field delimiter) is shown here:

```
"1","a string","100.20"
"2","a string containing a , comma","102.20"
"3","a string containing a \" quote","102.20"
"4","a string containing a \", quote and comma","102.20"
```

If you specify `OPTIONALLY`, the `ENCLOSED BY` character is used only to enclose values from columns that have a string data type (such as `CHAR`, `BINARY`, `TEXT`, or `ENUM`):

```
1,"a string",100.20
2,"a string containing a , comma",102.20
3,"a string containing a \" quote",102.20
4,"a string containing a \", quote and comma",102.20
```

Occurrences of the `ENCLOSED BY` character within a field value are escaped by prefixing them with the `ESCAPED BY` character. Also, if you specify an empty `ESCAPED BY` value, it is possible to inadvertently generate output that cannot be read properly by `LOAD DATA INFILE`. For example, the preceding output just shown would appear as follows if the escape character is empty. Observe that the second

field in the fourth line contains a comma following the quote, which (erroneously) appears to terminate the field:

```
1,"a string",100.20
2,"a string containing a , comma",102.20
3,"a string containing a " quote",102.20
4,"a string containing a ", quote and comma",102.20
```

For input, the `ENCLOSED BY` character, if present, is stripped from the ends of field values. (This is true regardless of whether `OPTIONALLY` is specified; `OPTIONALLY` has no effect on input interpretation.) Occurrences of the `ENCLOSED BY` character preceded by the `ESCAPED BY` character are interpreted as part of the current field value.

If the field begins with the `ENCLOSED BY` character, instances of that character are recognized as terminating a field value only if followed by the field or line `TERMINATED BY` sequence. To avoid ambiguity, occurrences of the `ENCLOSED BY` character within a field value can be doubled and are interpreted as a single instance of the character. For example, if `ENCLOSED BY '"'` is specified, quotation marks are handled as shown here:

```
"The ""BIG"" boss"  -> The "BIG" boss
The "BIG" boss      -> The "BIG" boss
The ""BIG"" boss    -> The ""BIG"" boss
```

`FIELDS ESCAPED BY` controls how to read or write special characters:

- For input, if the `FIELDS ESCAPED BY` character is not empty, occurrences of that character are stripped and the following character is taken literally as part of a field value. Some two-character sequences that are exceptions, where the first character is the escape character. These sequences are shown in the following table (using `\` for the escape character). The rules for `NULL` handling are described later in this section.

Character	Escape Sequence
`\0`	An ASCII NUL (`X'00'`) character
`\b`	A backspace character
`\n`	A newline (linefeed) character
`\r`	A carriage return character
`\t`	A tab character.
`\Z`	ASCII 26 (Control+Z)
`\N`	NULL

  For more information about `\`-escape syntax, see Section 9.1.1, “String Literals”.

  If the `FIELDS ESCAPED BY` character is empty, escape-sequence interpretation does not occur.

- For output, if the `FIELDS ESCAPED BY` character is not empty, it is used to prefix the following characters on output:

  - The `FIELDS ESCAPED BY` character
  - The `FIELDS [OPTIONALLY] ENCLOSED BY` character
  - The first character of the `FIELDS TERMINATED BY` and `LINES TERMINATED BY` values
  - ASCII `0` (what is actually written following the escape character is ASCII `0`, not a zero-valued byte)

  If the `FIELDS ESCAPED BY` character is empty, no characters are escaped and `NULL` is output as `NULL`, not `\N`. It is probably not a good idea to specify an empty escape character, particularly if field values in your data contain any of the characters in the list just given.

In certain cases, field- and line-handling options interact:

- If `LINES TERMINATED BY` is an empty string and `FIELDS TERMINATED BY` is nonempty, lines are also terminated with `FIELDS TERMINATED BY`.
- If the `FIELDS TERMINATED BY` and `FIELDS ENCLOSED BY` values are both empty (`''`), a fixed-row (nondelimited) format is used. With fixed-row format, no delimiters are used between fields (but you can still have a line terminator). Instead, column values are read and written using a field width wide enough to hold all values in the field. For `TINYINT`, `SMALLINT`, `MEDIUMINT`, `INT`, and `BIGINT`, the field widths are 4, 6, 8, 11, and 20, respectively, no matter what the declared display width is.

  `LINES TERMINATED BY` is still used to separate lines. If a line does not contain all fields, the rest of the columns are set to their default values. If you do not have a line terminator, you should set this to `''`. In this case, the text file must contain all fields for each row.

  Fixed-row format also affects handling of `NULL` values, as described later.

  Note

  Fixed-size format does not work if you are using a multibyte character set.

Handling of `NULL` values varies according to the `FIELDS` and `LINES` options in use:

- For the default `FIELDS` and `LINES` values, `NULL` is written as a field value of `\N` for output, and a field value of `\N` is read as `NULL` for input (assuming that the `ESCAPED BY` character is `\`).
- If `FIELDS ENCLOSED BY` is not empty, a field containing the literal word `NULL` as its value is read as a `NULL` value. This differs from the word `NULL` enclosed within `FIELDS ENCLOSED BY` characters, which is read as the string `'NULL'`.
- If `FIELDS ESCAPED BY` is empty, `NULL` is written as the word `NULL`.
- With fixed-row format (which is used when `FIELDS TERMINATED BY` and `FIELDS ENCLOSED BY` are both empty), `NULL` is written as an empty string. This causes both `NULL` values and empty strings in the table to be indistinguishable when written to the file because both are written as empty strings. If you need to be able to tell the two apart when reading the file back in, you should not use fixed-row format.

An attempt to load `NULL` into a `NOT NULL` column causes assignment of the implicit default value for the column's data type and a warning, or an error in strict SQL mode. Implicit default values are discussed in Section 11.7, “Data Type Default Values”.

Some cases are not supported by `LOAD DATA INFILE`:

- Fixed-size rows (`FIELDS TERMINATED BY` and `FIELDS ENCLOSED BY` both empty) and `BLOB` or `TEXT` columns.
- If you specify one separator that is the same as or a prefix of another, `LOAD DATA INFILE` cannot interpret the input properly. For example, the following `FIELDS` clause would cause problems:

```
FIELDS TERMINATED BY '"' ENCLOSED BY '"'
```

- If `FIELDS ESCAPED BY` is empty, a field value that contains an occurrence of `FIELDS ENCLOSED BY` or `LINES TERMINATED BY` followed by the `FIELDS TERMINATED BY` value causes `LOAD DATA INFILE` to stop reading a field or line too early. This happens because `LOAD DATA INFILE` cannot properly determine where the field or line value ends.

The following example loads all columns of the `persondata` table:

```
LOAD DATA INFILE 'persondata.txt' INTO TABLE persondata;
```

By default, when no column list is provided at the end of the `LOAD DATA INFILE` statement, input lines are expected to contain a field for each table column. If you want to load only some of a table's columns, specify a column list:

```
LOAD DATA INFILE 'persondata.txt' INTO TABLE persondata (col1,col2,...);
```

You must also specify a column list if the order of the fields in the input file differs from the order of the columns in the table. Otherwise, MySQL cannot tell how to match input fields with table columns.

The column list can contain either column names or user variables. With user variables, the `SET` clause enables you to perform transformations on their values before assigning the result to columns.

User variables in the `SET` clause can be used in several ways. The following example uses the first input column directly for the value of `t1.column1`, and assigns the second input column to a user variable that is subjected to a division operation before being used for the value of `t1.column2`:

```
LOAD DATA INFILE 'file.txt'
  INTO TABLE t1
  (column1, @var1)
  SET column2 = @var1/100;
```

The `SET` clause can be used to supply values not derived from the input file. The following statement sets `column3` to the current date and time:

```
LOAD DATA INFILE 'file.txt'
  INTO TABLE t1
  (column1, column2)
  SET column3 = CURRENT_TIMESTAMP;
```

You can also discard an input value by assigning it to a user variable and not assigning the variable to a table column:

```
LOAD DATA INFILE 'file.txt'
  INTO TABLE t1
  (column1, @dummy, column2, @dummy, column3);
```

Use of the column/variable list and `SET` clause is subject to the following restrictions:

- Assignments in the `SET` clause should have only column names on the left hand side of assignment operators.
- You can use subqueries in the right hand side of `SET` assignments. A subquery that returns a value to be assigned to a column may be a scalar subquery only. Also, you cannot use a subquery to select from the table that is being loaded.
- Lines ignored by an `IGNORE` clause are not processed for the column/variable list or `SET` clause.
- User variables cannot be used when loading data with fixed-row format because user variables do not have a display width.

When processing an input line, `LOAD DATA` splits it into fields and uses the values according to the column/variable list and the `SET` clause, if they are present. Then the resulting row is inserted into the table. If there are `BEFORE INSERT` or `AFTER INSERT` triggers for the table, they are activated before or after inserting the row, respectively.

If an input line has too many fields, the extra fields are ignored and the number of warnings is incremented.

If an input line has too few fields, the table columns for which input fields are missing are set to their default values. Default value assignment is described in Section 11.7, “Data Type Default Values”.

An empty field value is interpreted different from a missing field:

- For string types, the column is set to the empty string.
- For numeric types, the column is set to `0`.
- For date and time types, the column is set to the appropriate “zero” value for the type. See Section 11.3, “Date and Time Types”.

These are the same values that result if you assign an empty string explicitly to a string, numeric, or date or time type explicitly in an `INSERT` or `UPDATE` statement.

Treatment of empty or incorrect field values differs from that just described if the SQL mode is set to a restrictive value. For example, if `sql_mode` is set to `TRADITIONAL`, conversion of an empty value or a value such as `'x'` for a numeric column results in an error, not conversion to 0. (With `LOCAL` or `IGNORE`, warnings occur rather than errors, even with a restrictive `sql_mode` value, and the row is inserted using the same closest-value behavior used for nonrestrictive SQL modes. This occurs because the server has no way to stop transmission of the file in the middle of the operation.)

`TIMESTAMP` columns are set to the current date and time only if there is a `NULL` value for the column (that is, `\N`) and the column is not declared to permit `NULL` values, or if the `TIMESTAMP` column's default value is the current timestamp and it is omitted from the field list when a field list is specified.

`LOAD DATA INFILE` regards all input as strings, so you cannot use numeric values for `ENUM` or `SET` columns the way you can with `INSERT` statements. All `ENUM` and `SET` values must be specified as strings.

`BIT` values cannot be loaded directly using binary notation (for example, `b'011010'`). To work around this, use the `SET` clause to strip off the leading `b'` and trailing `'` and perform a base-2 to base-10 conversion so that MySQL loads the values into the `BIT` column properly:

```
shell> cat /tmp/bit_test.txt
b'10'
b'1111111'
shell> mysql test
mysql> LOAD DATA INFILE '/tmp/bit_test.txt'
       INTO TABLE bit_test (@var1)
       SET b = CAST(CONV(MID(@var1, 3, LENGTH(@var1)-3), 2, 10) AS UNSIGNED);
Query OK, 2 rows affected (0.00 sec)
Records: 2  Deleted: 0  Skipped: 0  Warnings: 0

mysql> SELECT BIN(b+0) FROM bit_test;
+----------+
| BIN(b+0) |
+----------+
| 10       |
| 1111111  |
+----------+
2 rows in set (0.00 sec)
```

For `BIT` values in `0b` binary notation (for example, `0b011010`), use this `SET` clause instead to strip off the leading `0b`:

```
SET b = CAST(CONV(MID(@var1, 3, LENGTH(@var1)-2), 2, 10) AS UNSIGNED)
```

On Unix, if you need `LOAD DATA` to read from a pipe, you can use the following technique (the example loads a listing of the `/` directory into the table `db1.t1`):

```
mkfifo /mysql/data/db1/ls.dat
chmod 666 /mysql/data/db1/ls.dat
find / -ls > /mysql/data/db1/ls.dat &
mysql -e "LOAD DATA INFILE 'ls.dat' INTO TABLE t1" db1
```

Here you must run the command that generates the data to be loaded and the `mysql` commands either on separate terminals, or run the data generation process in the background (as shown in the preceding example). If you do not do this, the pipe will block until data is read by the `mysql` process.

When the `LOAD DATA INFILE` statement finishes, it returns an information string in the following format:

```
Records: 1  Deleted: 0  Skipped: 0  Warnings: 0
```

Warnings occur under the same circumstances as when values are inserted using the `INSERT` statement (see Section 13.2.6, “INSERT Syntax”), except that `LOAD DATA INFILE` also generates warnings when there are too few or too many fields in the input row.

You can use `SHOW WARNINGS` to get a list of the first `max_error_count` warnings as information about what went wrong. See Section 13.7.5.40, “SHOW WARNINGS Syntax”.

If you are using the C API, you can get information about the statement by calling the `mysql_info()` function. See Section 27.7.7.36, “mysql_info()”.

13.2.8 LOAD XML Syntax

```
LOAD XML [LOW_PRIORITY | CONCURRENT] [LOCAL] INFILE 'file_name'
    [REPLACE | IGNORE]
    INTO TABLE [db_name.]tbl_name
    [CHARACTER SET charset_name]
    [ROWS IDENTIFIED BY '<tagname>']
    [IGNORE number {LINES | ROWS}]
    [(field_name_or_user_var,...)]
    [SET col_name = expr,...]
```

The `LOAD XML` statement reads data from an XML file into a table. The *`file_name`* must be given as a literal string. The *`tagname`* in the optional `ROWS IDENTIFIED BY` clause must also be given as a literal string, and must be surrounded by angle brackets (`<` and `>`).

`LOAD XML` acts as the complement of running the `mysql` client in XML output mode (that is, starting the client with the `--xml` option). To write data from a table to an XML file, you can invoke the `mysql` client with the `--xml` and `-e` options from the system shell, as shown here:

```
shell> mysql --xml -e 'SELECT * FROM mydb.mytable' > file.xml
```

To read the file back into a table, use `LOAD XML INFILE`. By default, the `<row>` element is considered to be the equivalent of a database table row; this can be changed using the `ROWS IDENTIFIED BY` clause.

This statement supports three different XML formats:

- Column names as attributes and column values as attribute values:

  ```
  <row column1="value1" column2="value2" .../>
  ```

- Column names as tags and column values as the content of these tags:

  ```
  <row>
    <column1>value1</column1>
  ```

```
  <column2>value2</column2>
</row>
```

- Column names are the `name` attributes of `<field>` tags, and values are the contents of these tags:

```
<row>
  <field name='column1'>value1</field>
  <field name='column2'>value2</field>
</row>
```

  This is the format used by other MySQL tools, such as `mysqldump`.

All three formats can be used in the same XML file; the import routine automatically detects the format for each row and interprets it correctly. Tags are matched based on the tag or attribute name and the column name.

The following clauses work essentially the same way for `LOAD XML` as they do for `LOAD DATA`:

- `LOW_PRIORITY` or `CONCURRENT`
- `LOCAL`
- `REPLACE` or `IGNORE`
- `CHARACTER SET`
- `SET`

See Section 13.2.7, “LOAD DATA INFILE Syntax”, for more information about these clauses.

`(field_name_or_user_var, ...)` is a comma-separated list of one or more XML fields or user variables. The name of a user variable used for this purpose must match the name of a field from the XML file, prefixed with `@`. You can use field names to select only desired fields. User variables can be employed to store the corresponding field values for subsequent re-use.

The `IGNORE` *`number`* `LINES` or `IGNORE` *`number`* `ROWS` clause causes the first *`number`* rows in the XML file to be skipped. It is analogous to the `LOAD DATA` statement's `IGNORE ... LINES` clause.

Suppose that we have a table named `person`, created as shown here:

```
USE test;

CREATE TABLE person (
    person_id INT NOT NULL PRIMARY KEY,
    fname VARCHAR(40) NULL,
    lname VARCHAR(40) NULL,
    created TIMESTAMP
);
```

Suppose further that this table is initially empty.

Now suppose that we have a simple XML file `person.xml`, whose contents are as shown here:

```
<list>
  <person person_id="1" fname="Kapek" lname="Sainnouine"/>
  <person person_id="2" fname="Sajon" lname="Rondela"/>
  <person person_id="3"><fname>Likame</fname><lname>Örrtmons</lname></person>
  <person person_id="4"><fname>Slar</fname><lname>Manlanth</lname></person>
  <person><field name="person_id">5</field><field name="fname">Stoma</field>
    <field name="lname">Milu</field></person>
  <person><field name="person_id">6</field><field name="fname">Nirtam</field>
    <field name="lname">Sklöd</field></person>
  <person person_id="7"><fname>Sungam</fname><lname>Dulbåd</lname></person>
```

```
  <person person_id="8" fname="Sraref" lname="Encmelt"/>
</list>
```

Each of the permissible XML formats discussed previously is represented in this example file.

To import the data in `person.xml` into the `person` table, you can use this statement:

```
mysql> LOAD XML LOCAL INFILE 'person.xml'
    ->   INTO TABLE person
    ->   ROWS IDENTIFIED BY '<person>';

Query OK, 8 rows affected (0.00 sec)
Records: 8  Deleted: 0  Skipped: 0  Warnings: 0
```

Here, we assume that `person.xml` is located in the MySQL data directory. If the file cannot be found, the following error results:

```
ERROR 2 (HY000): File '/person.xml' not found (Errcode: 2)
```

The `ROWS IDENTIFIED BY '<person>'` clause means that each `<person>` element in the XML file is considered equivalent to a row in the table into which the data is to be imported. In this case, this is the `person` table in the `test` database.

As can be seen by the response from the server, 8 rows were imported into the `test.person` table. This can be verified by a simple `SELECT` statement:

```
mysql> SELECT * FROM person;
+-----------+--------+------------+---------------------+
| person_id | fname  | lname      | created             |
+-----------+--------+------------+---------------------+
|         1 | Kapek  | Sainnouine | 2007-07-13 16:18:47 |
|         2 | Sajon  | Rondela    | 2007-07-13 16:18:47 |
|         3 | Likame | Örrtmons   | 2007-07-13 16:18:47 |
|         4 | Slar   | Manlanth   | 2007-07-13 16:18:47 |
|         5 | Stoma  | Nilu       | 2007-07-13 16:18:47 |
|         6 | Nirtam | Sklöd      | 2007-07-13 16:18:47 |
|         7 | Sungam | Dulbåd     | 2007-07-13 16:18:47 |
|         8 | Sreraf | Encmelt    | 2007-07-13 16:18:47 |
+-----------+--------+------------+---------------------+
8 rows in set (0.00 sec)
```

This shows, as stated earlier in this section, that any or all of the 3 permitted XML formats may appear in a single file and be read in using `LOAD XML`.

The inverse of the import operation just shown—that is, dumping MySQL table data into an XML file—can be accomplished using the `mysql` client from the system shell, as shown here:

```
shell> mysql --xml -e "SELECT * FROM test.person" > person-dump.xml
shell> cat person-dump.xml
<?xml version="1.0"?>

<resultset statement="SELECT * FROM test.person" xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance">
  <row>
 <field name="person_id">1</field>
 <field name="fname">Kapek</field>
 <field name="lname">Sainnouine</field>
  </row>

  <row>
 <field name="person_id">2</field>
 <field name="fname">Sajon</field>
 <field name="lname">Rondela</field>
  </row>
```

```
  <row>
 <field name="person_id">3</field>
 <field name="fname">Likema</field>
 <field name="lname">Örrtmons</field>
  </row>

  <row>
 <field name="person_id">4</field>
 <field name="fname">Slar</field>
 <field name="lname">Manlanth</field>
  </row>

  <row>
 <field name="person_id">5</field>
 <field name="fname">Stoma</field>
 <field name="lname">Nilu</field>
  </row>

  <row>
 <field name="person_id">6</field>
 <field name="fname">Nirtam</field>
 <field name="lname">Sklöd</field>
  </row>

  <row>
 <field name="person_id">7</field>
 <field name="fname">Sungam</field>
 <field name="lname">Dulbåd</field>
  </row>

  <row>
 <field name="person_id">8</field>
 <field name="fname">Sreraf</field>
 <field name="lname">Encmelt</field>
  </row>
</resultset>
```

Note

The `--xml` option causes the `mysql` client to use XML formatting for its output; the `-e` option causes the client to execute the SQL statement immediately following the option. See Section 4.5.1, “`mysql` — The MySQL Command-Line Tool”.

You can verify that the dump is valid by creating a copy of the `person` table and importing the dump file into the new table, like this:

```
mysql> USE test;
mysql> CREATE TABLE person2 LIKE person;
Query OK, 0 rows affected (0.00 sec)

mysql> LOAD XML LOCAL INFILE 'person-dump.xml'
    ->   INTO TABLE person2;
Query OK, 8 rows affected (0.01 sec)
Records: 8  Deleted: 0  Skipped: 0  Warnings: 0

mysql> SELECT * FROM person2;
+-----------+--------+------------+---------------------+
| person_id | fname  | lname      | created             |
+-----------+--------+------------+---------------------+
|         1 | Kapek  | Sainnouine | 2007-07-13 16:18:47 |
|         2 | Sajon  | Rondela    | 2007-07-13 16:18:47 |
|         3 | Likema | Örrtmons   | 2007-07-13 16:18:47 |
|         4 | Slar   | Manlanth   | 2007-07-13 16:18:47 |
|         5 | Stoma  | Nilu       | 2007-07-13 16:18:47 |
|         6 | Nirtam | Sklöd      | 2007-07-13 16:18:47 |
|         7 | Sungam | Dulbåd     | 2007-07-13 16:18:47 |
|         8 | Sreraf | Encmelt    | 2007-07-13 16:18:47 |
+-----------+--------+------------+---------------------+
8 rows in set (0.00 sec)
```

There is no requirement that every field in the XML file be matched with a column in the corresponding table. Fields which have no corresponding columns are skipped. You can see this by first emptying the `person2` table and dropping the `created` column, then using the same `LOAD XML` statement we just employed previously, like this:

```
mysql> TRUNCATE person2;
Query OK, 8 rows affected (0.26 sec)

mysql> ALTER TABLE person2 DROP COLUMN created;
Query OK, 0 rows affected (0.52 sec)
Records: 0  Duplicates: 0  Warnings: 0

mysql> SHOW CREATE TABLE person2\G
*************************** 1. row ***************************
       Table: person2
Create Table: CREATE TABLE `person2` (
  `person_id` int(11) NOT NULL,
  `fname` varchar(40) DEFAULT NULL,
  `lname` varchar(40) DEFAULT NULL,
  PRIMARY KEY (`person_id`)
) ENGINE=InnoDB DEFAULT CHARSET=utf8
1 row in set (0.00 sec)

mysql> LOAD XML LOCAL INFILE 'person-dump.xml'
    ->   INTO TABLE person2;
Query OK, 8 rows affected (0.01 sec)
Records: 8  Deleted: 0  Skipped: 0  Warnings: 0

mysql> SELECT * FROM person2;
+-----------+--------+------------+
| person_id | fname  | lname      |
+-----------+--------+------------+
|         1 | Kapek  | Sainnouine |
|         2 | Sajon  | Rondela    |
|         3 | Likema | Örrtmons   |
|         4 | Slar   | Manlanth   |
|         5 | Stoma  | Nilu       |
|         6 | Nirtam | Sklöd      |
|         7 | Sungam | Dulbåd     |
|         8 | Sreraf | Encmelt    |
+-----------+--------+------------+
8 rows in set (0.00 sec)
```

The order in which the fields are given within each row of the XML file does not affect the operation of `LOAD XML`; the field order can vary from row to row, and is not required to be in the same order as the corresponding columns in the table.

As mentioned previously, you can use a `(field_name_or_user_var, ...)` list of one or more XML fields (to select desired fields only) or user variables (to store the corresponding field values for later use). User variables can be especially useful when you want to insert data from an XML file into table columns whose names do not match those of the XML fields. To see how this works, we first create a table named `individual` whose structure matches that of the `person` table, but whose columns are named differently:

```
mysql> CREATE TABLE individual (
    ->     individual_id INT NOT NULL PRIMARY KEY,
    ->     name1 VARCHAR(40) NULL,
    ->     name2 VARCHAR(40) NULL,
    ->     made TIMESTAMP
    -> );
Query OK, 0 rows affected (0.42 sec)
```

In this case, you cannot simply load the XML file directly into the table, because the field and column names do not match:

```
mysql> LOAD XML INFILE '../bin/person-dump.xml' INTO TABLE test.individual;
```

```
ERROR 1263 (22004): Column set to default value; NULL supplied to NOT NULL column 'individual_id' at ro
```

This happens because the MySQL server looks for field names matching the column names of the target table. You can work around this problem by selecting the field values into user variables, then setting the target table's columns equal to the values of those variables using `SET`. You can perform both of these operations in a single statement, as shown here:

```
mysql> LOAD XML INFILE '../bin/person-dump.xml'
    ->     INTO TABLE test.individual (@person_id, @fname, @lname, @created)
    ->     SET individual_id=@person_id, name1=@fname, name2=@lname, made=@created;
Query OK, 8 rows affected (0.05 sec)
Records: 8  Deleted: 0  Skipped: 0  Warnings: 0

mysql> SELECT * FROM individual;
+---------------+--------+------------+---------------------+
| individual_id | name1  | name2      | made                |
+---------------+--------+------------+---------------------+
|             1 | Kapek  | Sainnouine | 2007-07-13 16:18:47 |
|             2 | Sajon  | Rondela    | 2007-07-13 16:18:47 |
|             3 | Likema | Örrtmons   | 2007-07-13 16:18:47 |
|             4 | Slar   | Manlanth   | 2007-07-13 16:18:47 |
|             5 | Stoma  | Nilu       | 2007-07-13 16:18:47 |
|             6 | Nirtam | Sklöd      | 2007-07-13 16:18:47 |
|             7 | Sungam | Dulbåd     | 2007-07-13 16:18:47 |
|             8 | Srraf  | Encmelt    | 2007-07-13 16:18:47 |
+---------------+--------+------------+---------------------+
8 rows in set (0.00 sec)
```

The names of the user variables *must* match those of the corresponding fields from the XML file, with the addition of the required `@` prefix to indicate that they are variables. The user variables need not be listed or assigned in the same order as the corresponding fields.

Using a `ROWS IDENTIFIED BY '<tagname>'` clause, it is possible to import data from the same XML file into database tables with different definitions. For this example, suppose that you have a file named `address.xml` which contains the following XML:

```
<?xml version="1.0"?>

<list>
  <person person_id="1">
    <fname>Robert</fname>
    <lname>Jones</lname>
    <address address_id="1" street="Mill Creek Road" zip="45365" city="Sidney"/>
    <address address_id="2" street="Main Street" zip="28681" city="Taylorsville"/>
  </person>

  <person person_id="2">
    <fname>Mary</fname>
    <lname>Smith</lname>
    <address address_id="3" street="River Road" zip="80239" city="Denver"/>
    <!-- <address address_id="4" street="North Street" zip="37920" city="Knoxville"/> -->
  </person>

</list>
```

You can again use the `test.person` table as defined previously in this section, after clearing all the existing records from the table and then showing its structure as shown here:

```
mysql< TRUNCATE person;
Query OK, 0 rows affected (0.04 sec)

mysql< SHOW CREATE TABLE person\G
*************************** 1. row ***************************
       Table: person
Create Table: CREATE TABLE `person` (
  `person_id` int(11) NOT NULL,
```

```
  `fname` varchar(40) DEFAULT NULL,
  `lname` varchar(40) DEFAULT NULL,
  `created` timestamp NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,
  PRIMARY KEY (`person_id`)
) ENGINE=MyISAM DEFAULT CHARSET=utf8mb4
1 row in set (0.00 sec)
```

Now create an `address` table in the `test` database using the following `CREATE TABLE` statement:

```
CREATE TABLE address (
    address_id INT NOT NULL PRIMARY KEY,
    person_id INT NULL,
    street VARCHAR(40) NULL,
    zip INT NULL,
    city VARCHAR(40) NULL,
    created TIMESTAMP
);
```

To import the data from the XML file into the `person` table, execute the following `LOAD XML` statement, which specifies that rows are to be specified by the `<person>` element, as shown here;

```
mysql> LOAD XML LOCAL INFILE 'address.xml'
    ->   INTO TABLE person
    ->   ROWS IDENTIFIED BY '<person>';
Query OK, 2 rows affected (0.00 sec)
Records: 2  Deleted: 0  Skipped: 0  Warnings: 0
```

You can verify that the records were imported using a `SELECT` statement:

```
mysql> SELECT * FROM person;
+-----------+--------+-------+---------------------+
| person_id | fname  | lname | created             |
+-----------+--------+-------+---------------------+
|         1 | Robert | Jones | 2007-07-24 17:37:06 |
|         2 | Mary   | Smith | 2007-07-24 17:37:06 |
+-----------+--------+-------+---------------------+
2 rows in set (0.00 sec)
```

Since the `<address>` elements in the XML file have no corresponding columns in the `person` table, they are skipped.

To import the data from the `<address>` elements into the `address` table, use the `LOAD XML` statement shown here:

```
mysql> LOAD XML LOCAL INFILE 'address.xml'
    ->   INTO TABLE address
    ->   ROWS IDENTIFIED BY '<address>';
Query OK, 3 rows affected (0.00 sec)
Records: 3  Deleted: 0  Skipped: 0  Warnings: 0
```

You can see that the data was imported using a `SELECT` statement such as this one:

```
mysql> SELECT * FROM address;
+------------+-----------+-----------------+-------+-------------+---------------------+
| address_id | person_id | street          | zip   | city        | created             |
+------------+-----------+-----------------+-------+-------------+---------------------+
|          1 |         1 | Mill Creek Road | 45365 | Sidney      | 2007-07-24 17:37:37 |
|          2 |         1 | Main Street     | 28681 | Taylorsville | 2007-07-24 17:37:37 |
|          3 |         2 | River Road      | 80239 | Denver      | 2007-07-24 17:37:37 |
+------------+-----------+-----------------+-------+-------------+---------------------+
3 rows in set (0.00 sec)
```

The data from the `<address>` element that is enclosed in XML comments is not imported. However, since there is a `person_id` column in the `address` table, the value of the `person_id` attribute from the parent `<person>` element for each `<address>` *is* imported into the `address` table.

Security Considerations. As with the `LOAD DATA` statement, the transfer of the XML file from the client host to the server host is initiated by the MySQL server. In theory, a patched server could be built that would tell the client program to transfer a file of the server's choosing rather than the file named by the client in the `LOAD XML` statement. Such a server could access any file on the client host to which the client user has read access.

In a Web environment, clients usually connect to MySQL from a Web server. A user that can run any command against the MySQL server can use `LOAD XML LOCAL` to read any files to which the Web server process has read access. In this environment, the client with respect to the MySQL server is actually the Web server, not the remote program being run by the user who connects to the Web server.

You can disable loading of XML files from clients by starting the server with `--local-infile=0` or `--local-infile=OFF`. This option can also be used when starting the `mysql` client to disable `LOAD XML` for the duration of the client session.

To prevent a client from loading XML files from the server, do not grant the `FILE` privilege to the corresponding MySQL user account, or revoke this privilege if the client user account already has it.

Important

Revoking the `FILE` privilege (or not granting it in the first place) keeps the user only from executing the `LOAD XML INFILE` statement (as well as the `LOAD_FILE()` function; it does *not* prevent the user from executing `LOAD XML LOCAL INFILE`. To disallow this statement, you must start the server or the client with `--local-infile=OFF`.

In other words, the `FILE` privilege affects only whether the client can read files on the server; it has no bearing on whether the client can read files on the local file system.

For partitioned tables using storage engines that employ table locks, such as `MyISAM`, any locks caused by `LOAD XML` perform locks on all partitions of the table. This does not apply to tables using storage engines which employ row-level locking, such as `InnoDB`. For more information, see Partitioning and Locking.

13.2.9 REPLACE Syntax

```
REPLACE [LOW_PRIORITY | DELAYED]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    [(col_name,...)]
    {VALUES | VALUE} ({expr | DEFAULT},...),(...),...
```

Or:

```
REPLACE [LOW_PRIORITY | DELAYED]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    SET col_name={expr | DEFAULT}, ...
```

Or:

```
REPLACE [LOW_PRIORITY | DELAYED]
    [INTO] tbl_name
    [PARTITION (partition_name,...)]
    [(col_name,...)]
    SELECT ...
```

`REPLACE` works exactly like `INSERT`, except that if an old row in the table has the same value as a new row for a `PRIMARY KEY` or a `UNIQUE` index, the old row is deleted before the new row is inserted. See Section 13.2.6, “INSERT Syntax”.

`REPLACE` is a MySQL extension to the SQL standard. It either inserts, or *deletes* and inserts. For another MySQL extension to standard SQL—that either inserts or *updates*—see Section 13.2.6.3, “INSERT ... ON DUPLICATE KEY UPDATE Syntax”.

`DELAYED` inserts and replaces were deprecated in MySQL 5.6.6. In MySQL 8.0, `DELAYED` is not supported. The server recognizes but ignores the `DELAYED` keyword, handles the replace as a nondelayed replace, and generates an `ER_WARN_LEGACY_SYNTAX_CONVERTED` warning. (“REPLACE DELAYED is no longer supported. The statement was converted to REPLACE.”) The `DELAYED` keyword will be removed in a future release.

Note

`REPLACE` makes sense only if a table has a `PRIMARY KEY` or `UNIQUE` index. Otherwise, it becomes equivalent to `INSERT`, because there is no index to be used to determine whether a new row duplicates another.

Values for all columns are taken from the values specified in the `REPLACE` statement. Any missing columns are set to their default values, just as happens for `INSERT`. You cannot refer to values from the current row and use them in the new row. If you use an assignment such as `SET col_name = col_name + 1`, the reference to the column name on the right hand side is treated as `DEFAULT(col_name)`, so the assignment is equivalent to `SET col_name = DEFAULT(col_name) + 1`.

If a generated column is replaced explicitly, the only permitted value is `DEFAULT`. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

To use `REPLACE`, you must have both the `INSERT` and `DELETE` privileges for the table.

`REPLACE` supports explicit partition selection using the `PARTITION` keyword with a comma-separated list of names of partitions, subpartitions, or both. As with `INSERT`, if it is not possible to insert the new row into any of these partitions or subpartitions, the `REPLACE` statement fails with the error `Found a row not matching the given partition set`. See Section 22.5, “Partition Selection”, for more information.

The `REPLACE` statement returns a count to indicate the number of rows affected. This is the sum of the rows deleted and inserted. If the count is 1 for a single-row `REPLACE`, a row was inserted and no rows were deleted. If the count is greater than 1, one or more old rows were deleted before the new row was inserted. It is possible for a single row to replace more than one old row if the table contains multiple unique indexes and the new row duplicates values for different old rows in different unique indexes.

The affected-rows count makes it easy to determine whether `REPLACE` only added a row or whether it also replaced any rows: Check whether the count is 1 (added) or greater (replaced).

If you are using the C API, the affected-rows count can be obtained using the `mysql_affected_rows()` function.

You cannot replace into a table and select from the same table in a subquery.

MySQL uses the following algorithm for `REPLACE` (and `LOAD DATA ... REPLACE`):

1. Try to insert the new row into the table
2. While the insertion fails because a duplicate-key error occurs for a primary key or unique index:
   a. Delete from the table the conflicting row that has the duplicate key value
   b. Try again to insert the new row into the table

It is possible that in the case of a duplicate-key error, a storage engine may perform the `REPLACE` as an update rather than a delete plus insert, but the semantics are the same. There are no user-visible

effects other than a possible difference in how the storage engine increments `Handler_xxx` status variables.

Because the results of `REPLACE ... SELECT` statements depend on the ordering of rows from the `SELECT` and this order cannot always be guaranteed, it is possible when logging these statements for the master and the slave to diverge. For this reason, `REPLACE ... SELECT` statements are flagged as unsafe for statement-based replication. With this change, such statements produce a warning in the log when using the `STATEMENT` binary logging mode, and are logged using the row-based format when using `MIXED` mode. See also Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

When modifying an existing table that is not partitioned to accommodate partitioning, or, when modifying the partitioning of an already partitioned table, you may consider altering the table's primary key (see Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”). You should be aware that, if you do this, the results of `REPLACE` statements may be affected, just as they would be if you modified the primary key of a nonpartitioned table. Consider the table created by the following `CREATE TABLE` statement:

```
CREATE TABLE test (
  id INT UNSIGNED NOT NULL AUTO_INCREMENT,
  data VARCHAR(64) DEFAULT NULL,
  ts TIMESTAMP NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,
  PRIMARY KEY (id)
);
```

When we create this table and run the statements shown in the mysql client, the result is as follows:

```
mysql> REPLACE INTO test VALUES (1, 'Old', '2014-08-20 18:47:00');
Query OK, 1 row affected (0.04 sec)

mysql> REPLACE INTO test VALUES (1, 'New', '2014-08-20 18:47:42');
Query OK, 2 rows affected (0.04 sec)

mysql> SELECT * FROM test;
+----+------+---------------------+
| id | data | ts                  |
+----+------+---------------------+
|  1 | New  | 2014-08-20 18:47:42 |
+----+------+---------------------+
1 row in set (0.00 sec)
```

Now we create a second table almost identical to the first, except that the primary key now covers 2 columns, as shown here (emphasized text):

```
CREATE TABLE test2 (
  id INT UNSIGNED NOT NULL AUTO_INCREMENT,
  data VARCHAR(64) DEFAULT NULL,
  ts TIMESTAMP NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,
  PRIMARY KEY (id, ts)
);
```

When we run on `test2` the same two `REPLACE` statements as we did on the original `test` table, we obtain a different result:

```
mysql> REPLACE INTO test2 VALUES (1, 'Old', '2014-08-20 18:47:00');
Query OK, 1 row affected (0.05 sec)

mysql> REPLACE INTO test2 VALUES (1, 'New', '2014-08-20 18:47:42');
Query OK, 1 row affected (0.06 sec)

mysql> SELECT * FROM test2;
+----+------+---------------------+
| id | data | ts                  |
```

```
+----+------+---------------------+
|  1 | Old  | 2014-08-20 18:47:00 |
|  1 | New  | 2014-08-20 18:47:42 |
+----+------+---------------------+
2 rows in set (0.00 sec)
```

This is due to the fact that, when run on `test2`, both the `id` and `ts` column values must match those of an existing row for the row to be replaced; otherwise, a row is inserted.

In MySQL 8.0, a `REPLACE` statement affecting a partitioned table using a storage engine such as `MyISAM` that employs table-level locks locks only those partitions containing rows that match the `REPLACE` statement's `WHERE` clause, as long as none of the table's partitioning columns are updated; otherwise the entire table is locked. (For storage engines such as `InnoDB` that employ row-level locking, no locking of partitions takes place.) For more information, see Partitioning and Locking.

13.2.10 SELECT Syntax

```
SELECT
    [ALL | DISTINCT | DISTINCTROW ]
      [HIGH_PRIORITY]
      [STRAIGHT_JOIN]
      [SQL_SMALL_RESULT] [SQL_BIG_RESULT] [SQL_BUFFER_RESULT]
      [SQL_CACHE | SQL_NO_CACHE] [SQL_CALC_FOUND_ROWS]
    select_expr [, select_expr ...]
    [FROM table_references
      [PARTITION partition_list]
    [WHERE where_condition]
    [GROUP BY {col_name | expr | position}
      [ASC | DESC], ... [WITH ROLLUP]]
    [HAVING where_condition]
    [ORDER BY {col_name | expr | position}
      [ASC | DESC], ...]
    [LIMIT {[offset,] row_count | row_count OFFSET offset}]
    [INTO OUTFILE 'file_name'
        [CHARACTER SET charset_name]
        export_options
      | INTO DUMPFILE 'file_name'
      | INTO var_name [, var_name]]
    [FOR UPDATE | LOCK IN SHARE MODE]]
    [FOR {UPDATE | SHARE} [OF tbl_name [, tbl_name] ...] [NOWAIT | SKIP LOCKED]
      | LOCK IN SHARE MODE]]
```

`SELECT` is used to retrieve rows selected from one or more tables, and can include `UNION` statements and subqueries. See Section 13.2.10.3, “UNION Syntax”, and Section 13.2.11, “Subquery Syntax”. A `SELECT` statement can start with a `WITH` clause to define common table expressions accessible within the `SELECT`. See Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”.

The most commonly used clauses of `SELECT` statements are these:

- Each *`select_expr`* indicates a column that you want to retrieve. There must be at least one *`select_expr`*.
- *`table_references`* indicates the table or tables from which to retrieve rows. Its syntax is described in Section 13.2.10.2, “JOIN Syntax”.
- `SELECT` supports explicit partition selection using the `PARTITION` with a list of partitions or subpartitions (or both) following the name of the table in a *`table_reference`* (see Section 13.2.10.2, “JOIN Syntax”). In this case, rows are selected only from the partitions listed, and any other partitions of the table are ignored. For more information and examples, see Section 22.5, “Partition Selection”.

  `SELECT ... PARTITION` from tables using storage engines such as `MyISAM` that perform table-level locks (and thus partition locks) lock only the partitions or subpartitions named by the `PARTITION` option.

See Partitioning and Locking, for more information.

- The WHERE clause, if given, indicates the condition or conditions that rows must satisfy to be selected. *where_condition* is an expression that evaluates to true for each row to be selected. The statement selects all rows if there is no WHERE clause.

  In the WHERE expression, you can use any of the functions and operators that MySQL supports, except for aggregate (summary) functions. See Section 9.5, “Expression Syntax”, and Chapter 12, *Functions and Operators*.

SELECT can also be used to retrieve rows computed without reference to any table.

For example:

```
mysql> SELECT 1 + 1;
        -> 2
```

You are permitted to specify DUAL as a dummy table name in situations where no tables are referenced:

```
mysql> SELECT 1 + 1 FROM DUAL;
        -> 2
```

DUAL is purely for the convenience of people who require that all SELECT statements should have FROM and possibly other clauses. MySQL may ignore the clauses. MySQL does not require FROM DUAL if no tables are referenced.

In general, clauses used must be given in exactly the order shown in the syntax description. For example, a HAVING clause must come after any GROUP BY clause and before any ORDER BY clause. The exception is that the INTO clause can appear either as shown in the syntax description or immediately following the *select_expr* list. For more information about INTO, see Section 13.2.10.1, “SELECT ... INTO Syntax”.

The list of *select_expr* terms comprises the select list that indicates which columns to retrieve. Terms specify a column or expression or can use *-shorthand:

- A select list consisting only of a single unqualified * can be used as shorthand to select all columns from all tables:

  ```
  SELECT * FROM t1 INNER JOIN t2 ...
  ```

- *tbl_name*.* can be used as a qualified shorthand to select all columns from the named table:

  ```
  SELECT t1.*, t2.* FROM t1 INNER JOIN t2 ...
  ```

- Use of an unqualified * with other items in the select list may produce a parse error. To avoid this problem, use a qualified *tbl_name*.* reference

  ```
  SELECT AVG(score), t1.* FROM t1 ...
  ```

The following list provides additional information about other SELECT clauses:

- A *select_expr* can be given an alias using AS *alias_name*. The alias is used as the expression's column name and can be used in GROUP BY, ORDER BY, or HAVING clauses. For example:

  ```
  ```

```
SELECT CONCAT(last_name,', ',first_name) AS full_name
  FROM mytable ORDER BY full_name;
```

The AS keyword is optional when aliasing a *select_expr* with an identifier. The preceding example could have been written like this:

```
SELECT CONCAT(last_name,', ',first_name) full_name
  FROM mytable ORDER BY full_name;
```

However, because the AS is optional, a subtle problem can occur if you forget the comma between two *select_expr* expressions: MySQL interprets the second as an alias name. For example, in the following statement, columnb is treated as an alias name:

```
SELECT columna columnb FROM mytable;
```

For this reason, it is good practice to be in the habit of using AS explicitly when specifying column aliases.

It is not permissible to refer to a column alias in a WHERE clause, because the column value might not yet be determined when the WHERE clause is executed. See Section B.5.4.4, “Problems with Column Aliases”.

- The FROM *table_references* clause indicates the table or tables from which to retrieve rows. If you name more than one table, you are performing a join. For information on join syntax, see Section 13.2.10.2, “JOIN Syntax”. For each table specified, you can optionally specify an alias.

```
tbl_name [[AS] alias] [index_hint]
```

The use of index hints provides the optimizer with information about how to choose indexes during query processing. For a description of the syntax for specifying these hints, see Section 8.9.4, “Index Hints”.

You can use SET max_seeks_for_key=*value* as an alternative way to force MySQL to prefer key scans instead of table scans. See Section 5.1.5, “Server System Variables”.

- You can refer to a table within the default database as *tbl_name*, or as *db_name*.*tbl_name* to specify a database explicitly. You can refer to a column as *col_name*, *tbl_name*.*col_name*, or *db_name*.*tbl_name*.*col_name*. You need not specify a *tbl_name* or *db_name*.*tbl_name* prefix for a column reference unless the reference would be ambiguous. See Section 9.2.1, “Identifier Qualifiers”, for examples of ambiguity that require the more explicit column reference forms.

- A table reference can be aliased using *tbl_name* AS *alias_name* or *tbl_name* *alias_name*:

```
SELECT t1.name, t2.salary FROM employee AS t1, info AS t2
  WHERE t1.name = t2.name;

SELECT t1.name, t2.salary FROM employee t1, info t2
  WHERE t1.name = t2.name;
```

- Columns selected for output can be referred to in ORDER BY and GROUP BY clauses using column names, column aliases, or column positions. Column positions are integers and begin with 1:

```
SELECT college, region, seed FROM tournament
  ORDER BY region, seed;

SELECT college, region AS r, seed AS s FROM tournament
  ORDER BY r, s;

SELECT college, region, seed FROM tournament
  ORDER BY 2, 3;
```

To sort in reverse order, add the `DESC` (descending) keyword to the name of the column in the `ORDER BY` clause that you are sorting by. The default is ascending order; this can be specified explicitly using the `ASC` keyword.

If `ORDER BY` occurs within a subquery and also is applied in the outer query, the outermost `ORDER BY` takes precedence. For example, results for the following statement are sorted in descending order, not ascending order:

```
(SELECT ... ORDER BY a) ORDER BY a DESC;
```

Use of column positions is deprecated because the syntax has been removed from the SQL standard.

- If you use `GROUP BY`, output rows are sorted according to the `GROUP BY` columns as if you had an `ORDER BY` for the same columns. To avoid the overhead of sorting that `GROUP BY` produces, add `ORDER BY NULL`:

```
SELECT a, COUNT(b) FROM test_table GROUP BY a ORDER BY NULL;
```

Previously, relying on implicit `GROUP BY` sorting was deprecated but `GROUP BY` did sort by default (that is, in the absence of `ASC` or `DESC` designators). In MySQL 8.0, `GROUP BY` no longer sorts by default, so query results may differ from previous MySQL versions. To produce a given sort order, use explicit `ASC` or `DESC` designators for `GROUP BY` columns or provide an `ORDER BY` clause.

- When you use `ORDER BY` or `GROUP BY` to sort a column in a `SELECT`, the server sorts values using only the initial number of bytes indicated by the `max_sort_length` system variable.
- MySQL extends the `GROUP BY` clause so that you can also specify `ASC` and `DESC` after columns named in the clause:

```
SELECT a, COUNT(b) FROM test_table GROUP BY a DESC;
```

- MySQL extends the use of `GROUP BY` to permit selecting fields that are not mentioned in the `GROUP BY` clause. If you are not getting the results that you expect from your query, please read the description of `GROUP BY` found in Section 12.18, “Aggregate (GROUP BY) Functions”.
- `GROUP BY` permits a `WITH ROLLUP` modifier. See Section 12.18.2, “GROUP BY Modifiers”.
- The `HAVING` clause is applied nearly last, just before items are sent to the client, with no optimization. (`LIMIT` is applied after `HAVING`.)

  The SQL standard requires that `HAVING` must reference only columns in the `GROUP BY` clause or columns used in aggregate functions. However, MySQL supports an extension to this behavior, and permits `HAVING` to refer to columns in the `SELECT` list and columns in outer subqueries as well.

  If the `HAVING` clause refers to a column that is ambiguous, a warning occurs. In the following statement, `col2` is ambiguous because it is used as both an alias and a column name:

```
SELECT COUNT(col1) AS col2 FROM t GROUP BY col2 HAVING col2 = 2;
```

  Preference is given to standard SQL behavior, so if a `HAVING` column name is used both in `GROUP BY` and as an aliased column in the output column list, preference is given to the column in the `GROUP BY` column.

- Do not use `HAVING` for items that should be in the `WHERE` clause. For example, do not write the following:

```
SELECT col_name FROM tbl_name HAVING col_name > 0;
```

Write this instead:

```
SELECT col_name FROM tbl_name WHERE col_name > 0;
```

- The HAVING clause can refer to aggregate functions, which the WHERE clause cannot:

```
SELECT user, MAX(salary) FROM users
  GROUP BY user HAVING MAX(salary) > 10;
```

(This did not work in some older versions of MySQL.)

- MySQL permits duplicate column names. That is, there can be more than one *select_expr* with the same name. This is an extension to standard SQL. Because MySQL also permits GROUP BY and HAVING to refer to *select_expr* values, this can result in an ambiguity:

```
SELECT 12 AS a, a FROM t GROUP BY a;
```

In that statement, both columns have the name a. To ensure that the correct column is used for grouping, use different names for each *select_expr*.

- MySQL resolves unqualified column or alias references in ORDER BY clauses by searching in the *select_expr* values, then in the columns of the tables in the FROM clause. For GROUP BY or HAVING clauses, it searches the FROM clause before searching in the *select_expr* values. (For GROUP BY and HAVING, this differs from the pre-MySQL 5.0 behavior that used the same rules as for ORDER BY.)

- The LIMIT clause can be used to constrain the number of rows returned by the SELECT statement. LIMIT takes one or two numeric arguments, which must both be nonnegative integer constants, with these exceptions:

  - Within prepared statements, LIMIT parameters can be specified using ? placeholder markers.

  - Within stored programs, LIMIT parameters can be specified using integer-valued routine parameters or local variables.

  With two arguments, the first argument specifies the offset of the first row to return, and the second specifies the maximum number of rows to return. The offset of the initial row is 0 (not 1):

```
SELECT * FROM tbl LIMIT 5,10;  # Retrieve rows 6-15
```

To retrieve all rows from a certain offset up to the end of the result set, you can use some large number for the second parameter. This statement retrieves all rows from the 96th row to the last:

```
SELECT * FROM tbl LIMIT 95,18446744073709551615;
```

With one argument, the value specifies the number of rows to return from the beginning of the result set:

```
SELECT * FROM tbl LIMIT 5;     # Retrieve first 5 rows
```

In other words, LIMIT *row_count* is equivalent to LIMIT 0, *row_count*.

For prepared statements, you can use placeholders. The following statements will return one row from the tbl table:

```
SET @a=1;
PREPARE STMT FROM 'SELECT * FROM tbl LIMIT ?';
```

```
EXECUTE STMT USING @a;
```

The following statements will return the second to sixth row from the `tbl` table:

```
SET @skip=1; SET @numrows=5;
PREPARE STMT FROM 'SELECT * FROM tbl LIMIT ?, ?';
EXECUTE STMT USING @skip, @numrows;
```

For compatibility with PostgreSQL, MySQL also supports the `LIMIT row_count OFFSET offset` syntax.

If `LIMIT` occurs within a subquery and also is applied in the outer query, the outermost `LIMIT` takes precedence. For example, the following statement produces two rows, not one:

```
(SELECT ... LIMIT 1) LIMIT 2;
```

- The `SELECT ... INTO` form of `SELECT` enables the query result to be written to a file or stored in variables. For more information, see Section 13.2.10.1, “SELECT ... INTO Syntax”.

- If you use `FOR UPDATE` with a storage engine that uses page or row locks, rows examined by the query are write-locked until the end of the current transaction.

  You cannot use `FOR UPDATE` as part of the `SELECT` in a statement such as `CREATE TABLE new_table SELECT ... FROM old_table ...`. (If you attempt to do so, the statement is rejected with the error `Can't update table 'old_table' while 'new_table' is being created`.)

  `FOR SHARE` and `LOCK IN SHARE MODE` set shared locks that permit other transactions to read the examined rows but not to update or delete them. `FOR SHARE` and `LOCK IN SHARE MODE` are equivalent. However, `FOR SHARE`, like `FOR UPDATE`, supports `NOWAIT`, `SKIP LOCKED`, and `OF tbl_name` options. `FOR SHARE` is a replacement for `LOCK IN SHARE MODE`, but `LOCK IN SHARE MODE` remains available for backward compatibility.

  `NOWAIT` causes a `FOR UPDATE` or `FOR SHARE` query to execute immediately, returning an error if a row lock cannot be obtained due to a lock held by another transaction.

  `SKIP LOCKED` causes a `FOR UPDATE` or `FOR SHARE` query to execute immediately, excluding rows from the result set that are locked by another transaction.

  `NOWAIT` and `SKIP LOCKED` options are unsafe for statement-based replication.

  Note

  Queries that skip locked rows return an inconsistent view of the data. `SKIP LOCKED` is therefore not suitable for general transactional work. However, it may be used to avoid lock contention when multiple sessions access the same queue-like table.

  `OF tbl_name` applies `FOR UPDATE` and `FOR SHARE` queries to named tables. For example:

```
SELECT * FROM t1, t2 FOR SHARE OF t1 FOR UPDATE OF t2;
```

  All tables referenced by the query block are locked when `OF tbl_name` is omitted. Consequently, using a locking clause without `OF tbl_name` in combination with another locking clause returns an error. Specifying the same table in multiple locking clauses returns an error. If an alias is specified as the table name in the `SELECT` statement, a locking clause may only use the alias. If the `SELECT` statement does not specify an alias explicitly, the locking clause may only specify the actual table name.

For more information about FOR UPDATE and FOR SHARE, see Section 15.5.2.4, "Locking Reads". For additional information about NOWAIT and SKIP LOCKED options, see Locking Read Concurrency with NOWAIT and SKIP LOCKED.

Following the SELECT keyword, you can use a number of modifiers that affect the operation of the statement. HIGH_PRIORITY, STRAIGHT_JOIN, and modifiers beginning with SQL_ are MySQL extensions to standard SQL.

- The ALL and DISTINCT modifiers specify whether duplicate rows should be returned. ALL (the default) specifies that all matching rows should be returned, including duplicates. DISTINCT specifies removal of duplicate rows from the result set. It is an error to specify both modifiers. DISTINCTROW is a synonym for DISTINCT.

- HIGH_PRIORITY gives the SELECT higher priority than a statement that updates a table. You should use this only for queries that are very fast and must be done at once. A SELECT HIGH_PRIORITY query that is issued while the table is locked for reading runs even if there is an update statement waiting for the table to be free. This affects only storage engines that use only table-level locking (such as MyISAM, MEMORY, and MERGE).

  HIGH_PRIORITY cannot be used with SELECT statements that are part of a UNION.

- STRAIGHT_JOIN forces the optimizer to join the tables in the order in which they are listed in the FROM clause. You can use this to speed up a query if the optimizer joins the tables in nonoptimal order. STRAIGHT_JOIN also can be used in the *table_references* list. See Section 13.2.10.2, "JOIN Syntax".

  STRAIGHT_JOIN does not apply to any table that the optimizer treats as a const or system table. Such a table produces a single row, is read during the optimization phase of query execution, and references to its columns are replaced with the appropriate column values before query execution proceeds. These tables will appear first in the query plan displayed by EXPLAIN. See Section 8.8.1, "Optimizing Queries with EXPLAIN". This exception may not apply to const or system tables that are used on the NULL-complemented side of an outer join (that is, the right-side table of a LEFT JOIN or the left-side table of a RIGHT JOIN.

- SQL_BIG_RESULT or SQL_SMALL_RESULT can be used with GROUP BY or DISTINCT to tell the optimizer that the result set has many rows or is small, respectively. For SQL_BIG_RESULT, MySQL directly uses disk-based temporary tables if needed, and prefers sorting to using a temporary table with a key on the GROUP BY elements. For SQL_SMALL_RESULT, MySQL uses fast temporary tables to store the resulting table instead of using sorting. This should not normally be needed.

- SQL_BUFFER_RESULT forces the result to be put into a temporary table. This helps MySQL free the table locks early and helps in cases where it takes a long time to send the result set to the client. This modifier can be used only for top-level SELECT statements, not for subqueries or following UNION.

- SQL_CALC_FOUND_ROWS tells MySQL to calculate how many rows there would be in the result set, disregarding any LIMIT clause. The number of rows can then be retrieved with SELECT FOUND_ROWS(). See Section 12.14, "Information Functions".

- The SQL_CACHE and SQL_NO_CACHE modifiers affect caching of query results in the query cache (see Section 8.10.3, "The MySQL Query Cache"). SQL_CACHE tells MySQL to store the result in the query cache if it is cacheable and the value of the query_cache_type system variable is 2 or DEMAND. With SQL_NO_CACHE, the server does not use the query cache. It neither checks the query cache to see whether the result is already cached, nor does it cache the query result.

  These two modifiers are mutually exclusive and an error occurs if they are both specified. Also, these modifiers are not permitted in subqueries (including subqueries in the FROM clause), and SELECT statements in unions other than the first SELECT.

For views, SQL_NO_CACHE applies if it appears in any SELECT in the query. For a cacheable query, SQL_CACHE applies if it appears in the first SELECT of a view referred to by the query.

In MySQL 8.0, a SELECT from a partitioned table using a storage engine such as MyISAM that employs table-level locks locks only those partitions containing rows that match the SELECT statement's WHERE clause. (This does not occur with storage engines such as InnoDB that employ row-level locking.) For more information, see Partitioning and Locking.

13.2.10.1 SELECT ... INTO Syntax

The SELECT ... INTO form of SELECT enables a query result to be stored in variables or written to a file:

- SELECT ... INTO *var_list* selects column values and stores them into variables.
- SELECT ... INTO OUTFILE writes the selected rows to a file. Column and line terminators can be specified to produce a specific output format.
- SELECT ... INTO DUMPFILE writes a single row to a file without any formatting.

The SELECT syntax description (see Section 13.2.10, “SELECT Syntax”) shows the INTO clause near the end of the statement. It is also possible to use INTO immediately following the *select_expr* list.

An INTO clause should not be used in a nested SELECT because such a SELECT must return its result to the outer context.

The INTO clause can name a list of one or more variables, which can be user-defined variables, stored procedure or function parameters, or stored program local variables. (Within a prepared SELECT ... INTO OUTFILE statement, only user-defined variables are permitted;see Section 13.6.4.2, “Local Variable Scope and Resolution”.)

The selected values are assigned to the variables. The number of variables must match the number of columns. The query should return a single row. If the query returns no rows, a warning with error code 1329 occurs (No data), and the variable values remain unchanged. If the query returns multiple rows, error 1172 occurs (Result consisted of more than one row). If it is possible that the statement may retrieve multiple rows, you can use LIMIT 1 to limit the result set to a single row.

```
SELECT id, data INTO @x, @y FROM test.t1 LIMIT 1;
```

User variable names are not case sensitive. See Section 9.4, “User-Defined Variables”.

The SELECT ... INTO OUTFILE '*file_name*' form of SELECT writes the selected rows to a file. The file is created on the server host, so you must have the FILE privilege to use this syntax. *file_name* cannot be an existing file, which among other things prevents files such as /etc/passwd and database tables from being destroyed. The character_set_filesystem system variable controls the interpretation of the file name.

The SELECT ... INTO OUTFILE statement is intended primarily to let you very quickly dump a table to a text file on the server machine. If you want to create the resulting file on some other host than the server host, you normally cannot use SELECT ... INTO OUTFILE since there is no way to write a path to the file relative to the server host's file system.

However, if the MySQL client software is installed on the remote machine, you can instead use a client command such as mysql -e "SELECT ..." > *file_name* to generate the file on the client host.

It is also possible to create the resulting file on a different host other than the server host, if the location of the file on the remote host can be accessed using a network-mapped path on the server's file system. In this case, the presence of mysql (or some other MySQL client program) is not required on the target host.

`SELECT ... INTO OUTFILE` is the complement of `LOAD DATA INFILE`. Column values are written converted to the character set specified in the `CHARACTER SET` clause. If no such clause is present, values are dumped using the `binary` character set. In effect, there is no character set conversion. If a result set contains columns in several character sets, the output data file will as well and you may not be able to reload the file correctly.

The syntax for the *`export_options`* part of the statement consists of the same `FIELDS` and `LINES` clauses that are used with the `LOAD DATA INFILE` statement. See Section 13.2.7, "LOAD DATA INFILE Syntax", for information about the `FIELDS` and `LINES` clauses, including their default values and permissible values.

`FIELDS ESCAPED BY` controls how to write special characters. If the `FIELDS ESCAPED BY` character is not empty, it is used when necessary to avoid ambiguity as a prefix that precedes following characters on output:

- The `FIELDS ESCAPED BY` character
- The `FIELDS [OPTIONALLY] ENCLOSED BY` character
- The first character of the `FIELDS TERMINATED BY` and `LINES TERMINATED BY` values
- ASCII `NUL` (the zero-valued byte; what is actually written following the escape character is ASCII `0`, not a zero-valued byte)

The `FIELDS TERMINATED BY`, `ENCLOSED BY`, `ESCAPED BY`, or `LINES TERMINATED BY` characters *must* be escaped so that you can read the file back in reliably. ASCII `NUL` is escaped to make it easier to view with some pagers.

The resulting file does not have to conform to SQL syntax, so nothing else need be escaped.

If the `FIELDS ESCAPED BY` character is empty, no characters are escaped and `NULL` is output as `NULL`, not `\N`. It is probably not a good idea to specify an empty escape character, particularly if field values in your data contain any of the characters in the list just given.

Here is an example that produces a file in the comma-separated values (CSV) format used by many programs:

```
SELECT a,b,a+b INTO OUTFILE '/tmp/result.txt'
  FIELDS TERMINATED BY ',' OPTIONALLY ENCLOSED BY '"'
  LINES TERMINATED BY '\n'
  FROM test_table;
```

If you use `INTO DUMPFILE` instead of `INTO OUTFILE`, MySQL writes only one row into the file, without any column or line termination and without performing any escape processing. This is useful if you want to store a `BLOB` value in a file.

Note

Any file created by `INTO OUTFILE` or `INTO DUMPFILE` is writable by all users on the server host. The reason for this is that the MySQL server cannot create a file that is owned by anyone other than the user under whose account it is running. (You should *never* run `mysqld` as `root` for this and other reasons.) The file thus must be world-writable so that you can manipulate its contents.

If the `secure_file_priv` system variable is set to a nonempty directory name, the file to be written must be located in that directory.

In the context of `SELECT ... INTO` statements that occur as part of events executed by the Event Scheduler, diagnostics messages (not only errors, but also warnings) are written to the error log, and, on Windows, to the application event log. For additional information, see Section 23.4.5, "Event Scheduler Status".

13.2.10.2 JOIN Syntax

MySQL supports the following `JOIN` syntax for the *`table_references`* part of `SELECT` statements and multiple-table `DELETE` and `UPDATE` statements:

```
table_references:
    escaped_table_reference [, escaped_table_reference] ...

escaped_table_reference:
    table_reference
  | { OJ table_reference }

table_reference:
    table_factor
  | join_table

table_factor:
    tbl_name [PARTITION (partition_names)]
        [[AS] alias] [index_hint_list]
  | table_subquery [AS] alias [(col_list)]
  | ( table_references )

join_table:
    table_reference [INNER | CROSS] JOIN table_factor [join_condition]
  | table_reference STRAIGHT_JOIN table_factor
  | table_reference STRAIGHT_JOIN table_factor ON conditional_expr
  | table_reference {LEFT|RIGHT} [OUTER] JOIN table_reference join_condition
  | table_reference NATURAL [INNER | {LEFT|RIGHT} [OUTER]] JOIN table_factor

join_condition:
    ON conditional_expr
  | USING (column_list)

index_hint_list:
    index_hint [, index_hint] ...

index_hint:
    USE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] ([index_list])
  | IGNORE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] (index_list)
  | FORCE {INDEX|KEY}
      [FOR {JOIN|ORDER BY|GROUP BY}] (index_list)

index_list:
    index_name [, index_name] ...
```

A table reference is also known as a join expression.

A table reference (when it refers to a partitioned table) may contain a `PARTITION` option, including a comma-separated list of partitions, subpartitions, or both. This option follows the name of the table and precedes any alias declaration. The effect of this option is that rows are selected only from the listed partitions or subpartitions. Any partitions or subpartitions not named in the list are ignored. For more information, see Section 22.5, “Partition Selection”.

The syntax of *`table_factor`* is extended in MySQL in comparison with standard SQL. The standard accepts only *`table_reference`*, not a list of them inside a pair of parentheses.

This is a conservative extension if each comma in a list of *`table_reference`* items is considered as equivalent to an inner join. For example:

```
SELECT * FROM t1 LEFT JOIN (t2, t3, t4)
                 ON (t2.a = t1.a AND t3.b = t1.b AND t4.c = t1.c)
```

is equivalent to:

```
SELECT * FROM t1 LEFT JOIN (t2 CROSS JOIN t3 CROSS JOIN t4)
```

```
                ON (t2.a = t1.a AND t3.b = t1.b AND t4.c = t1.c)
```

In MySQL, `JOIN`, `CROSS JOIN`, and `INNER JOIN` are syntactic equivalents (they can replace each other). In standard SQL, they are not equivalent. `INNER JOIN` is used with an `ON` clause, `CROSS JOIN` is used otherwise.

In general, parentheses can be ignored in join expressions containing only inner join operations. MySQL also supports nested joins. See Section 8.2.1.7, “Nested Join Optimization”.

Index hints can be specified to affect how the MySQL optimizer makes use of indexes. For more information, see Section 8.9.4, “Index Hints”. Optimizer hints and the `optimizer_switch` system variable are other ways to influence optimizer use of indexes. See Section 8.9.2, “Optimizer Hints”, and Section 8.9.3, “Switchable Optimizations”.

The following list describes general factors to take into account when writing joins:

- A table reference can be aliased using *`tbl_name AS alias_name`* or *`tbl_name alias_name`*:

  ```
  SELECT t1.name, t2.salary
    FROM employee AS t1 INNER JOIN info AS t2 ON t1.name = t2.name;

  SELECT t1.name, t2.salary
    FROM employee t1 INNER JOIN info t2 ON t1.name = t2.name;
  ```

- A *`table_subquery`* is also known as a derived table or subquery in the `FROM` clause. See Section 13.2.11.8, “Derived Tables (Subqueries in the FROM Clause)”. Such subqueries *must* include an alias to give the subquery result a table name, and may optionally include a list of table column names in parentheses. A trivial example follows:

  ```
  SELECT * FROM (SELECT 1, 2, 3) AS t1;
  ```

- `INNER JOIN` and `,` (comma) are semantically equivalent in the absence of a join condition: both produce a Cartesian product between the specified tables (that is, each and every row in the first table is joined to each and every row in the second table).

  However, the precedence of the comma operator is less than that of `INNER JOIN`, `CROSS JOIN`, `LEFT JOIN`, and so on. If you mix comma joins with the other join types when there is a join condition, an error of the form `Unknown column 'col_name' in 'on clause'` may occur. Information about dealing with this problem is given later in this section.

- The *`conditional_expr`* used with `ON` is any conditional expression of the form that can be used in a `WHERE` clause. Generally, the `ON` clause serves for conditions that specify how to join tables, and the `WHERE` clause restricts which rows to include in the result set.

- If there is no matching row for the right table in the `ON` or `USING` part in a `LEFT JOIN`, a row with all columns set to `NULL` is used for the right table. You can use this fact to find rows in a table that have no counterpart in another table:

  ```
  SELECT left_tbl.*
    FROM left_tbl LEFT JOIN right_tbl ON left_tbl.id = right_tbl.id
    WHERE right_tbl.id IS NULL;
  ```

  This example finds all rows in `left_tbl` with an `id` value that is not present in `right_tbl` (that is, all rows in `left_tbl` with no corresponding row in `right_tbl`). See Section 8.2.1.8, “Left Join and Right Join Optimization”.

- The `USING(column_list)` clause names a list of columns that must exist in both tables. If tables `a` and `b` both contain columns `c1`, `c2`, and `c3`, the following join compares corresponding columns from the two tables:

  ```
  a LEFT JOIN b USING (c1, c2, c3)
  ```

- The `NATURAL [LEFT] JOIN` of two tables is defined to be semantically equivalent to an `INNER JOIN` or a `LEFT JOIN` with a `USING` clause that names all columns that exist in both tables.
- `RIGHT JOIN` works analogously to `LEFT JOIN`. To keep code portable across databases, it is recommended that you use `LEFT JOIN` instead of `RIGHT JOIN`.
- The `{ OJ ... }` syntax shown in the join syntax description exists only for compatibility with ODBC. The curly braces in the syntax should be written literally; they are not metasyntax as used elsewhere in syntax descriptions.

```
SELECT left_tbl.*
    FROM { OJ left_tbl LEFT OUTER JOIN right_tbl ON left_tbl.id = right_tbl.id }
    WHERE right_tbl.id IS NULL;
```

  You can use other types of joins within `{ OJ ... }`, such as `INNER JOIN` or `RIGHT OUTER JOIN`. This helps with compatibility with some third-party applications, but is not official ODBC syntax.
- `STRAIGHT_JOIN` is similar to `JOIN`, except that the left table is always read before the right table. This can be used for those (few) cases for which the join optimizer processes the tables in a suboptimal order.

Some join examples:

```
SELECT * FROM table1, table2;

SELECT * FROM table1 INNER JOIN table2 ON table1.id = table2.id;

SELECT * FROM table1 LEFT JOIN table2 ON table1.id = table2.id;

SELECT * FROM table1 LEFT JOIN table2 USING (id);

SELECT * FROM table1 LEFT JOIN table2 ON table1.id = table2.id
  LEFT JOIN table3 ON table2.id = table3.id;
```

Natural joins and joins with `USING`, including outer join variants, are processed according to the SQL:2003 standard:

- Redundant columns of a `NATURAL` join do not appear. Consider this set of statements:

```
CREATE TABLE t1 (i INT, j INT);
CREATE TABLE t2 (k INT, j INT);
INSERT INTO t1 VALUES(1, 1);
INSERT INTO t2 VALUES(1, 1);
SELECT * FROM t1 NATURAL JOIN t2;
SELECT * FROM t1 JOIN t2 USING (j);
```

  In the first `SELECT` statement, column `j` appears in both tables and thus becomes a join column, so, according to standard SQL, it should appear only once in the output, not twice. Similarly, in the second SELECT statement, column `j` is named in the `USING` clause and should appear only once in the output, not twice.

  Thus, the statements produce this output:

```
+------+------+------+
| j    | i    | k    |
+------+------+------+
|    1 |    1 |    1 |
+------+------+------+
+------+------+------+
| j    | i    | k    |
+------+------+------+
|    1 |    1 |    1 |
+------+------+------+
```

Redundant column elimination and column ordering occurs according to standard SQL, producing this display order:

- First, coalesced common columns of the two joined tables, in the order in which they occur in the first table
- Second, columns unique to the first table, in order in which they occur in that table
- Third, columns unique to the second table, in order in which they occur in that table

The single result column that replaces two common columns is defined using the coalesce operation. That is, for two `t1.a` and `t2.a` the resulting single join column `a` is defined as `a = COALESCE(t1.a, t2.a)`, where:

```
COALESCE(x, y) = (CASE WHEN x IS NOT NULL THEN x ELSE y END)
```

If the join operation is any other join, the result columns of the join consist of the concatenation of all columns of the joined tables.

A consequence of the definition of coalesced columns is that, for outer joins, the coalesced column contains the value of the non-`NULL` column if one of the two columns is always `NULL`. If neither or both columns are `NULL`, both common columns have the same value, so it doesn't matter which one is chosen as the value of the coalesced column. A simple way to interpret this is to consider that a coalesced column of an outer join is represented by the common column of the inner table of a `JOIN`. Suppose that the tables `t1(a, b)` and `t2(a, c)` have the following contents:

```
t1    t2
----  ----
1 x   2 z
2 y   3 w
```

Then, for this join, column `a` contains the values of `t1.a`:

```
mysql> SELECT * FROM t1 NATURAL LEFT JOIN t2;
+------+------+------+
| a    | b    | c    |
+------+------+------+
|    1 | x    | NULL |
|    2 | y    | z    |
+------+------+------+
```

By contrast, for this join, column `a` contains the values of `t2.a`.

```
mysql> SELECT * FROM t1 NATURAL RIGHT JOIN t2;
+------+------+------+
| a    | c    | b    |
+------+------+------+
|    2 | z    | y    |
|    3 | w    | NULL |
+------+------+------+
```

Compare those results to the otherwise equivalent queries with `JOIN ... ON`:

```
mysql> SELECT * FROM t1 LEFT JOIN t2 ON (t1.a = t2.a);
+------+------+------+------+
| a    | b    | a    | c    |
+------+------+------+------+
|    1 | x    | NULL | NULL |
|    2 | y    |    2 | z    |
+------+------+------+------+
```

```
mysql> SELECT * FROM t1 RIGHT JOIN t2 ON (t1.a = t2.a);
+------+------+------+------+
| a    | b    | a    | c    |
+------+------+------+------+
|    2 | y    |    2 | z    |
| NULL | NULL |    3 | w    |
+------+------+------+------+
```

- A `USING` clause can be rewritten as an `ON` clause that compares corresponding columns. However, although `USING` and `ON` are similar, they are not quite the same. Consider the following two queries:

  ```
  a LEFT JOIN b USING (c1, c2, c3)
  a LEFT JOIN b ON a.c1 = b.c1 AND a.c2 = b.c2 AND a.c3 = b.c3
  ```

  With respect to determining which rows satisfy the join condition, both joins are semantically identical.

  With respect to determining which columns to display for `SELECT *` expansion, the two joins are not semantically identical. The `USING` join selects the coalesced value of corresponding columns, whereas the `ON` join selects all columns from all tables. For the `USING` join, `SELECT *` selects these values:

  ```
  COALESCE(a.c1, b.c1), COALESCE(a.c2, b.c2), COALESCE(a.c3, b.c3)
  ```

  For the `ON` join, `SELECT *` selects these values:

  ```
  a.c1, a.c2, a.c3, b.c1, b.c2, b.c3
  ```

  With an inner join, `COALESCE(a.c1, b.c1)` is the same as either `a.c1` or `b.c1` because both columns will have the same value. With an outer join (such as `LEFT JOIN`), one of the two columns can be `NULL`. That column is omitted from the result.

- An `ON` clause can refer only to its operands.

  Example:

  ```
  CREATE TABLE t1 (i1 INT);
  CREATE TABLE t2 (i2 INT);
  CREATE TABLE t3 (i3 INT);
  SELECT * FROM t1 JOIN t2 ON (i1 = i3) JOIN t3;
  ```

  The statement fails with an `Unknown column 'i3' in 'on clause'` error because `i3` is a column in `t3`, which is not an operand of the `ON` clause. To enable the join to be processed, rewrite the statement as follows:

  ```
  SELECT * FROM t1 JOIN t2 JOIN t3 ON (i1 = i3);
  ```

- `JOIN` has higher precedence than the comma operator (`,`), so the join expression `t1, t2 JOIN t3` is interpreted as `(t1, (t2 JOIN t3))`, not as `((t1, t2) JOIN t3)`. This affects statements that use an `ON` clause because that clause can refer only to columns in the operands of the join, and the precedence affects interpretation of what those operands are.

  Example:

  ```
  CREATE TABLE t1 (i1 INT, j1 INT);
  CREATE TABLE t2 (i2 INT, j2 INT);
  CREATE TABLE t3 (i3 INT, j3 INT);
  INSERT INTO t1 VALUES(1, 1);
  INSERT INTO t2 VALUES(1, 1);
  ```

```
INSERT INTO t3 VALUES(1, 1);
SELECT * FROM t1, t2 JOIN t3 ON (t1.i1 = t3.i3);
```

The `JOIN` takes precedence over the comma operator, so the operands for the `ON` clause are `t2` and `t3`. Because `t1.i1` is not a column in either of the operands, the result is an `Unknown column 't1.i1' in 'on clause'` error.

To enable the join to be processed, use either of these strategies:

- Group the first two tables explicitly with parentheses so that the operands for the `ON` clause are `(t1, t2)` and `t3`:

  ```
  SELECT * FROM (t1, t2) JOIN t3 ON (t1.i1 = t3.i3);
  ```

- Avoid the use of the comma operator and use `JOIN` instead:

  ```
  SELECT * FROM t1 JOIN t2 JOIN t3 ON (t1.i1 = t3.i3);
  ```

The same precedence interpretation also applies to statements that mix the comma operator with `INNER JOIN`, `CROSS JOIN`, `LEFT JOIN`, and `RIGHT JOIN`, all of which have higher precedence than the comma operator.

- A MySQL extension compared to the SQL:2003 standard is that MySQL permits you to qualify the common (coalesced) columns of `NATURAL` or `USING` joins, whereas the standard disallows that.

13.2.10.3 UNION Syntax

```
SELECT ...
UNION [ALL | DISTINCT] SELECT ...
[UNION [ALL | DISTINCT] SELECT ...]
```

`UNION` is used to combine the result from multiple `SELECT` statements into a single result set.

The column names from the first `SELECT` statement are used as the column names for the results returned. Selected columns listed in corresponding positions of each `SELECT` statement should have the same data type. (For example, the first column selected by the first statement should have the same type as the first column selected by the other statements.)

If the data types of corresponding `SELECT` columns do not match, the types and lengths of the columns in the `UNION` result take into account the values retrieved by all of the `SELECT` statements. For example, consider the following:

```
mysql> SELECT REPEAT('a',1) UNION SELECT REPEAT('b',10);
+---------------+
| REPEAT('a',1) |
+---------------+
| a             |
| bbbbbbbbbb    |
+---------------+
```

The `SELECT` statements are normal select statements, but with the following restrictions:

- Only the last `SELECT` statement can use `INTO OUTFILE`. (However, the entire `UNION` result is written to the file.)

- `HIGH_PRIORITY` cannot be used with `SELECT` statements that are part of a `UNION`. If you specify it for the first `SELECT`, it has no effect. If you specify it for any subsequent `SELECT` statements, a syntax error results.

The default behavior for `UNION` is that duplicate rows are removed from the result. The optional `DISTINCT` keyword has no effect other than the default because it also specifies duplicate-row

removal. With the optional `ALL` keyword, duplicate-row removal does not occur and the result includes all matching rows from all the `SELECT` statements.

You can mix `UNION ALL` and `UNION DISTINCT` in the same query. Mixed `UNION` types are treated such that a `DISTINCT` union overrides any `ALL` union to its left. A `DISTINCT` union can be produced explicitly by using `UNION DISTINCT` or implicitly by using `UNION` with no following `DISTINCT` or `ALL` keyword.

To apply `ORDER BY` or `LIMIT` to an individual `SELECT`, place the clause inside the parentheses that enclose the `SELECT`:

```
(SELECT a FROM t1 WHERE a=10 AND B=1 ORDER BY a LIMIT 10)
UNION
(SELECT a FROM t2 WHERE a=11 AND B=2 ORDER BY a LIMIT 10);
```

However, use of `ORDER BY` for individual `SELECT` statements implies nothing about the order in which the rows appear in the final result because `UNION` by default produces an unordered set of rows. Therefore, the use of `ORDER BY` in this context is typically in conjunction with `LIMIT`, so that it is used to determine the subset of the selected rows to retrieve for the `SELECT`, even though it does not necessarily affect the order of those rows in the final `UNION` result. If `ORDER BY` appears without `LIMIT` in a `SELECT`, it is optimized away because it will have no effect anyway.

To use an `ORDER BY` or `LIMIT` clause to sort or limit the entire `UNION` result, parenthesize the individual `SELECT` statements and place the `ORDER BY` or `LIMIT` after the last one. The following example uses both clauses:

```
(SELECT a FROM t1 WHERE a=10 AND B=1)
UNION
(SELECT a FROM t2 WHERE a=11 AND B=2)
ORDER BY a LIMIT 10;
```

A statement without parentheses is equivalent to one parenthesized as just shown.

This kind of `ORDER BY` cannot use column references that include a table name (that is, names in *`tbl_name.col_name`* format). Instead, provide a column alias in the first `SELECT` statement and refer to the alias in the `ORDER BY`. (Alternatively, refer to the column in the `ORDER BY` using its column position. However, use of column positions is deprecated.)

Also, if a column to be sorted is aliased, the `ORDER BY` clause *must* refer to the alias, not the column name. The first of the following statements will work, but the second will fail with an `Unknown column 'a' in 'order clause'` error:

```
(SELECT a AS b FROM t) UNION (SELECT ...) ORDER BY b;
(SELECT a AS b FROM t) UNION (SELECT ...) ORDER BY a;
```

To cause rows in a `UNION` result to consist of the sets of rows retrieved by each `SELECT` one after the other, select an additional column in each `SELECT` to use as a sort column and add an `ORDER BY` following the last `SELECT`:

```
(SELECT 1 AS sort_col, col1a, col1b, ... FROM t1)
UNION
(SELECT 2, col2a, col2b, ... FROM t2) ORDER BY sort_col;
```

To additionally maintain sort order within individual `SELECT` results, add a secondary column to the `ORDER BY` clause:

```
(SELECT 1 AS sort_col, col1a, col1b, ... FROM t1)
UNION
(SELECT 2, col2a, col2b, ... FROM t2) ORDER BY sort_col, col1a;
```

Use of an additional column also enables you to determine which `SELECT` each row comes from. Extra columns can provide other identifying information as well, such as a string that indicates a table name.

`UNION` queries with an aggregate function in an `ORDER BY` clause are rejected with an `ER_AGGREGATE_ORDER_FOR_UNION` error. Example:

```
SELECT 1 AS foo UNION SELECT 2 ORDER BY MAX(1);
```

13.2.11 Subquery Syntax

A subquery is a `SELECT` statement within another statement.

All subquery forms and operations that the SQL standard requires are supported, as well as a few features that are MySQL-specific.

Here is an example of a subquery:

```
SELECT * FROM t1 WHERE column1 = (SELECT column1 FROM t2);
```

In this example, `SELECT * FROM t1 ...` is the *outer query* (or *outer statement*), and `(SELECT column1 FROM t2)` is the *subquery*. We say that the subquery is *nested* within the outer query, and in fact it is possible to nest subqueries within other subqueries, to a considerable depth. A subquery must always appear within parentheses.

The main advantages of subqueries are:

- They allow queries that are *structured* so that it is possible to isolate each part of a statement.
- They provide alternative ways to perform operations that would otherwise require complex joins and unions.
- Many people find subqueries more readable than complex joins or unions. Indeed, it was the innovation of subqueries that gave people the original idea of calling the early SQL “Structured Query Language.”

Here is an example statement that shows the major points about subquery syntax as specified by the SQL standard and supported in MySQL:

```
DELETE FROM t1
WHERE s11 > ANY
 (SELECT COUNT(*) /* no hint */ FROM t2
  WHERE NOT EXISTS
   (SELECT * FROM t3
    WHERE ROW(5*t2.s1,77)=
     (SELECT 50,11*s1 FROM t4 UNION SELECT 50,77 FROM
      (SELECT * FROM t5) AS t5)));
```

A subquery can return a scalar (a single value), a single row, a single column, or a table (one or more rows of one or more columns). These are called scalar, column, row, and table subqueries. Subqueries that return a particular kind of result often can be used only in certain contexts, as described in the following sections.

There are few restrictions on the type of statements in which subqueries can be used. A subquery can contain many of the keywords or clauses that an ordinary `SELECT` can contain: `DISTINCT`, `GROUP BY`, `ORDER BY`, `LIMIT`, joins, index hints, `UNION` constructs, comments, functions, and so on.

A subquery's outer statement can be any one of: `SELECT`, `INSERT`, `UPDATE`, `DELETE`, `SET`, or `DO`.

In MySQL, you cannot modify a table and select from the same table in a subquery. This applies to statements such as `DELETE`, `INSERT`, `REPLACE`, `UPDATE`, and (because subqueries can be used in the `SET` clause) `LOAD DATA INFILE`.

For information about how the optimizer handles subqueries, see Section 8.2.2, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions”. For a discussion of restrictions on subquery use, including performance issues for certain forms of subquery syntax, see Section C.4, “Restrictions on Subqueries”.

13.2.11.1 The Subquery as Scalar Operand

In its simplest form, a subquery is a scalar subquery that returns a single value. A scalar subquery is a simple operand, and you can use it almost anywhere a single column value or literal is legal, and you can expect it to have those characteristics that all operands have: a data type, a length, an indication that it can be `NULL`, and so on. For example:

```
CREATE TABLE t1 (s1 INT, s2 CHAR(5) NOT NULL);
INSERT INTO t1 VALUES(100, 'abcde');
SELECT (SELECT s2 FROM t1);
```

The subquery in this `SELECT` returns a single value (`'abcde'`) that has a data type of `CHAR`, a length of 5, a character set and collation equal to the defaults in effect at `CREATE TABLE` time, and an indication that the value in the column can be `NULL`. Nullability of the value selected by a scalar subquery is not copied because if the subquery result is empty, the result is `NULL`. For the subquery just shown, if `t1` were empty, the result would be `NULL` even though `s2` is `NOT NULL`.

There are a few contexts in which a scalar subquery cannot be used. If a statement permits only a literal value, you cannot use a subquery. For example, `LIMIT` requires literal integer arguments, and `LOAD DATA INFILE` requires a literal string file name. You cannot use subqueries to supply these values.

When you see examples in the following sections that contain the rather spartan construct `(SELECT column1 FROM t1)`, imagine that your own code contains much more diverse and complex constructions.

Suppose that we make two tables:

```
CREATE TABLE t1 (s1 INT);
INSERT INTO t1 VALUES (1);
CREATE TABLE t2 (s1 INT);
INSERT INTO t2 VALUES (2);
```

Then perform a `SELECT`:

```
SELECT (SELECT s1 FROM t2) FROM t1;
```

The result is `2` because there is a row in `t2` containing a column `s1` that has a value of `2`.

A scalar subquery can be part of an expression, but remember the parentheses, even if the subquery is an operand that provides an argument for a function. For example:

```
SELECT UPPER((SELECT s1 FROM t1)) FROM t2;
```

13.2.11.2 Comparisons Using Subqueries

The most common use of a subquery is in the form:

```
non_subquery_operand comparison_operator (subquery)
```

Where *`comparison_operator`* is one of these operators:

```
=  >  <  >=  <=  <>  !=  <=>
```

For example:

```
... WHERE 'a' = (SELECT column1 FROM t1)
```

MySQL also permits this construct:

```
non_subquery_operand LIKE (subquery)
```

At one time the only legal place for a subquery was on the right side of a comparison, and you might still find some old DBMSs that insist on this.

Here is an example of a common-form subquery comparison that you cannot do with a join. It finds all the rows in table `t1` for which the `column1` value is equal to a maximum value in table `t2`:

```
SELECT * FROM t1
  WHERE column1 = (SELECT MAX(column2) FROM t2);
```

Here is another example, which again is impossible with a join because it involves aggregating for one of the tables. It finds all rows in table `t1` containing a value that occurs twice in a given column:

```
SELECT * FROM t1 AS t
  WHERE 2 = (SELECT COUNT(*) FROM t1 WHERE t1.id = t.id);
```

For a comparison of the subquery to a scalar, the subquery must return a scalar. For a comparison of the subquery to a row constructor, the subquery must be a row subquery that returns a row with the same number of values as the row constructor. See Section 13.2.11.5, “Row Subqueries”.

13.2.11.3 Subqueries with ANY, IN, or SOME

Syntax:

```
operand comparison_operator ANY (subquery)
operand IN (subquery)
operand comparison_operator SOME (subquery)
```

Where *`comparison_operator`* is one of these operators:

```
=  >  <  >=  <=  <>  !=
```

The `ANY` keyword, which must follow a comparison operator, means “return `TRUE` if the comparison is `TRUE` for `ANY` of the values in the column that the subquery returns.” For example:

```
SELECT s1 FROM t1 WHERE s1 > ANY (SELECT s1 FROM t2);
```

Suppose that there is a row in table `t1` containing `(10)`. The expression is `TRUE` if table `t2` contains `(21,14,7)` because there is a value `7` in `t2` that is less than `10`. The expression is `FALSE` if table `t2` contains `(20,10)`, or if table `t2` is empty. The expression is *unknown* (that is, `NULL`) if table `t2` contains `(NULL,NULL,NULL)`.

When used with a subquery, the word `IN` is an alias for `= ANY`. Thus, these two statements are the same:

```
SELECT s1 FROM t1 WHERE s1 = ANY (SELECT s1 FROM t2);
SELECT s1 FROM t1 WHERE s1 IN    (SELECT s1 FROM t2);
```

`IN` and `= ANY` are not synonyms when used with an expression list. `IN` can take an expression list, but `= ANY` cannot. See Section 12.3.2, “Comparison Functions and Operators”.

`NOT IN` is not an alias for `<> ANY`, but for `<> ALL`. See Section 13.2.11.4, “Subqueries with ALL”.

The word `SOME` is an alias for `ANY`. Thus, these two statements are the same:

```
SELECT s1 FROM t1 WHERE s1 <> ANY  (SELECT s1 FROM t2);
SELECT s1 FROM t1 WHERE s1 <> SOME (SELECT s1 FROM t2);
```

Use of the word `SOME` is rare, but this example shows why it might be useful. To most people, the English phrase “a is not equal to any b” means “there is no b which is equal to a,” but that is not what is meant by the SQL syntax. The syntax means “there is some b to which a is not equal.” Using `<> SOME` instead helps ensure that everyone understands the true meaning of the query.

13.2.11.4 Subqueries with ALL

Syntax:

```
operand comparison_operator ALL (subquery)
```

The word `ALL`, which must follow a comparison operator, means “return `TRUE` if the comparison is `TRUE` for `ALL` of the values in the column that the subquery returns.” For example:

```
SELECT s1 FROM t1 WHERE s1 > ALL (SELECT s1 FROM t2);
```

Suppose that there is a row in table `t1` containing `(10)`. The expression is `TRUE` if table `t2` contains `(-5,0,+5)` because `10` is greater than all three values in `t2`. The expression is `FALSE` if table `t2` contains `(12,6,NULL,-100)` because there is a single value `12` in table `t2` that is greater than `10`. The expression is *unknown* (that is, `NULL`) if table `t2` contains `(0,NULL,1)`.

Finally, the expression is `TRUE` if table `t2` is empty. So, the following expression is `TRUE` when table `t2` is empty:

```
SELECT * FROM t1 WHERE 1 > ALL (SELECT s1 FROM t2);
```

But this expression is `NULL` when table `t2` is empty:

```
SELECT * FROM t1 WHERE 1 > (SELECT s1 FROM t2);
```

In addition, the following expression is `NULL` when table `t2` is empty:

```
SELECT * FROM t1 WHERE 1 > ALL (SELECT MAX(s1) FROM t2);
```

In general, *tables containing `NULL` values* and *empty tables* are “edge cases.” When writing subqueries, always consider whether you have taken those two possibilities into account.

`NOT IN` is an alias for `<> ALL`. Thus, these two statements are the same:

```
SELECT s1 FROM t1 WHERE s1 <> ALL (SELECT s1 FROM t2);
SELECT s1 FROM t1 WHERE s1 NOT IN (SELECT s1 FROM t2);
```

13.2.11.5 Row Subqueries

Scalar or column subqueries return a single value or a column of values. A *row subquery* is a subquery variant that returns a single row and can thus return more than one column value. Legal operators for row subquery comparisons are:

```
=  >  <  >=  <=  <>  !=  <=>
```

Here are two examples:

```
SELECT * FROM t1
  WHERE (col1,col2) = (SELECT col3, col4 FROM t2 WHERE id = 10);
SELECT * FROM t1
  WHERE ROW(col1,col2) = (SELECT col3, col4 FROM t2 WHERE id = 10);
```

For both queries, if the table `t2` contains a single row with `id = 10`, the subquery returns a single row. If this row has `col3` and `col4` values equal to the `col1` and `col2` values of any rows in `t1`, the `WHERE` expression is `TRUE` and each query returns those `t1` rows. If the `t2` row `col3` and `col4` values are not equal the `col1` and `col2` values of any `t1` row, the expression is `FALSE` and the query returns an empty result set. The expression is *unknown* (that is, `NULL`) if the subquery produces no rows. An error occurs if the subquery produces multiple rows because a row subquery can return at most one row.

For information about how each operator works for row comparisons, see Section 12.3.2, “Comparison Functions and Operators”.

The expressions `(1,2)` and `ROW(1,2)` are sometimes called *row constructors*. The two are equivalent. The row constructor and the row returned by the subquery must contain the same number of values.

A row constructor is used for comparisons with subqueries that return two or more columns. When a subquery returns a single column, this is regarded as a scalar value and not as a row, so a row constructor cannot be used with a subquery that does not return at least two columns. Thus, the following query fails with a syntax error:

```
SELECT * FROM t1 WHERE ROW(1) = (SELECT column1 FROM t2)
```

Row constructors are legal in other contexts. For example, the following two statements are semantically equivalent (and are handled in the same way by the optimizer):

```
SELECT * FROM t1 WHERE (column1,column2) = (1,1);
SELECT * FROM t1 WHERE column1 = 1 AND column2 = 1;
```

The following query answers the request, “find all rows in table `t1` that also exist in table `t2`”:

```
SELECT column1,column2,column3
  FROM t1
  WHERE (column1,column2,column3) IN
         (SELECT column1,column2,column3 FROM t2);
```

For more information about the optimizer and row constructors, see Section 8.2.1.18, “Row Constructor Expression Optimization”

13.2.11.6 Subqueries with EXISTS or NOT EXISTS

If a subquery returns any rows at all, `EXISTS subquery` is `TRUE`, and `NOT EXISTS subquery` is `FALSE`. For example:

```
SELECT column1 FROM t1 WHERE EXISTS (SELECT * FROM t2);
```

Traditionally, an `EXISTS` subquery starts with `SELECT *`, but it could begin with `SELECT 5` or `SELECT column1` or anything at all. MySQL ignores the `SELECT` list in such a subquery, so it makes no difference.

For the preceding example, if `t2` contains any rows, even rows with nothing but `NULL` values, the `EXISTS` condition is `TRUE`. This is actually an unlikely example because a `[NOT] EXISTS` subquery almost always contains correlations. Here are some more realistic examples:

- What kind of store is present in one or more cities?

```
SELECT DISTINCT store_type FROM stores
  WHERE EXISTS (SELECT * FROM cities_stores
                WHERE cities_stores.store_type = stores.store_type);
```

- What kind of store is present in no cities?

```
SELECT DISTINCT store_type FROM stores
  WHERE NOT EXISTS (SELECT * FROM cities_stores
                    WHERE cities_stores.store_type = stores.store_type);
```

- What kind of store is present in all cities?

```
SELECT DISTINCT store_type FROM stores s1
  WHERE NOT EXISTS (
    SELECT * FROM cities WHERE NOT EXISTS (
      SELECT * FROM cities_stores
       WHERE cities_stores.city = cities.city
       AND cities_stores.store_type = stores.store_type));
```

The last example is a double-nested `NOT EXISTS` query. That is, it has a `NOT EXISTS` clause within a `NOT EXISTS` clause. Formally, it answers the question “does a city exist with a store that is not in `Stores`”? But it is easier to say that a nested `NOT EXISTS` answers the question “is *x* `TRUE` for all *y*?”

13.2.11.7 Correlated Subqueries

A *correlated subquery* is a subquery that contains a reference to a table that also appears in the outer query. For example:

```
SELECT * FROM t1
  WHERE column1 = ANY (SELECT column1 FROM t2
                       WHERE t2.column2 = t1.column2);
```

Notice that the subquery contains a reference to a column of `t1`, even though the subquery's `FROM` clause does not mention a table `t1`. So, MySQL looks outside the subquery, and finds `t1` in the outer query.

Suppose that table `t1` contains a row where `column1 = 5` and `column2 = 6`; meanwhile, table `t2` contains a row where `column1 = 5` and `column2 = 7`. The simple expression `... WHERE column1 = ANY (SELECT column1 FROM t2)` would be `TRUE`, but in this example, the `WHERE` clause within the subquery is `FALSE` (because `(5,6)` is not equal to `(5,7)`), so the expression as a whole is `FALSE`.

Scoping rule: MySQL evaluates from inside to outside. For example:

```
SELECT column1 FROM t1 AS x
  WHERE x.column1 = (SELECT column1 FROM t2 AS x
    WHERE x.column1 = (SELECT column1 FROM t3
      WHERE x.column2 = t3.column1));
```

In this statement, `x.column2` must be a column in table `t2` because `SELECT column1 FROM t2 AS x ...` renames `t2`. It is not a column in table `t1` because `SELECT column1 FROM t1 ...` is an outer query that is *farther out*.

For subqueries in `HAVING` or `ORDER BY` clauses, MySQL also looks for column names in the outer select list.

For certain cases, a correlated subquery is optimized. For example:

```
val IN (SELECT key_val FROM tbl_name WHERE correlated_condition)
```

Otherwise, they are inefficient and likely to be slow. Rewriting the query as a join might improve performance.

Aggregate functions in correlated subqueries may contain outer references, provided the function contains nothing but outer references, and provided the function is not contained in another function or expression.

13.2.11.8 Derived Tables (Subqueries in the FROM Clause)

A derived table is a subquery in a `SELECT` statement `FROM` clause:

```
SELECT ... FROM (subquery) [AS] tbl_name ...
```

The `[AS] tbl_name` clause is mandatory because every table in a `FROM` clause must have a name. Any columns in the *`subquery`* select list must have unique names. Alternatively, *`tbl_name`* may be followed by a parenthesized list of names for the derived table columns:

```
SELECT ... FROM (subquery) [AS] tbl_name (col_list) ...
```

The number of names must be the same as the number of table columns.

For the sake of illustration, assume that you have this table:

```
CREATE TABLE t1 (s1 INT, s2 CHAR(5), s3 FLOAT);
```

Here is how to use a subquery in the `FROM` clause, using the example table:

```
INSERT INTO t1 VALUES (1,'1',1.0);
INSERT INTO t1 VALUES (2,'2',2.0);
SELECT sb1,sb2,sb3
  FROM (SELECT s1 AS sb1, s2 AS sb2, s3*2 AS sb3 FROM t1) AS sb
  WHERE sb1 > 1;
```

Result: `2, '2', 4.0`.

Here is another example: Suppose that you want to know the average of a set of sums for a grouped table. This does not work:

```
SELECT AVG(SUM(column1)) FROM t1 GROUP BY column1;
```

However, this query provides the desired information:

```
SELECT AVG(sum_column1)
  FROM (SELECT SUM(column1) AS sum_column1
        FROM t1 GROUP BY column1) AS t1;
```

Notice that the column name used within the subquery (`sum_column1`) is recognized in the outer query.

The column names for this derived table come from its select list:

```
mysql> SELECT * FROM (SELECT 1, 2, 3, 4) AS dt;
+---+---+---+---+
| 1 | 2 | 3 | 4 |
```

```
+---+---+---+---+
| 1 | 2 | 3 | 4 |
+---+---+---+---+
```

To provide column names, follow the derived table name with a parenthesized list of column names:

```
mysql> SELECT * FROM (SELECT 1, 2, 3, 4) AS dt (a, b, c, d);
+---+---+---+---+
| a | b | c | d |
+---+---+---+---+
| 1 | 2 | 3 | 4 |
+---+---+---+---+
```

Derived tables can return a scalar, column, row, or table.

Derived tables cannot be correlated subqueries, or contain outer references or references to other tables of the same `SELECT`.

The optimizer determines information about derived tables in such a way that materialization of them does not occur for `EXPLAIN`. See Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

It is possible under certain circumstances that using `EXPLAIN SELECT` will modify table data. This can occur if the outer query accesses any tables and an inner query invokes a stored function that changes one or more rows of a table. Suppose that there are two tables `t1` and `t2` in database `d1`, and a stored function `f1` that modifies `t2`, created as shown here:

```
CREATE DATABASE d1;
USE d1;
CREATE TABLE t1 (c1 INT);
CREATE TABLE t2 (c1 INT);
CREATE FUNCTION f1(p1 INT) RETURNS INT
  BEGIN
    INSERT INTO t2 VALUES (p1);
    RETURN p1;
  END;
```

Referencing the function directly in an `EXPLAIN SELECT` has no effect on `t2`, as shown here:

```
mysql> SELECT * FROM t2;
Empty set (0.02 sec)

mysql> EXPLAIN SELECT f1(5)\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: NULL
   partitions: NULL
         type: NULL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: NULL
     filtered: NULL
        Extra: No tables used
1 row in set (0.01 sec)

mysql> SELECT * FROM t2;
Empty set (0.01 sec)
```

This is because the `SELECT` statement did not reference any tables, as can be seen in the `table` and `Extra` columns of the output. This is also true of the following nested `SELECT`:

```
```

```
mysql> EXPLAIN SELECT NOW() AS a1, (SELECT f1(5)) AS a2\G
*************************** 1. row ***************************
           id: 1
  select_type: PRIMARY
        table: NULL
         type: NULL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: NULL
     filtered: NULL
        Extra: No tables used
1 row in set, 1 warning (0.00 sec)

mysql> SHOW WARNINGS;
+-------+------+-----------------------------------------+
| Level | Code | Message                                 |
+-------+------+-----------------------------------------+
| Note  | 1249 | Select 2 was reduced during optimization |
+-------+------+-----------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT * FROM t2;
Empty set (0.00 sec)
```

However, if the outer `SELECT` references any tables, the optimizer executes the statement in the subquery as well:

```
mysql> EXPLAIN SELECT * FROM t1 AS a1, (SELECT f1(5)) AS a2\G
*************************** 1. row ***************************
           id: 1
  select_type: PRIMARY
        table: <derived2>
   partitions: NULL
         type: system
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 1
     filtered: 100.00
        Extra: NULL
*************************** 2. row ***************************
           id: 1
  select_type: PRIMARY
        table: a1
   partitions: NULL
         type: ALL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 1
     filtered: 100.00
        Extra: NULL
*************************** 3. row ***************************
           id: 2
  select_type: DERIVED
        table: NULL
   partitions: NULL
         type: NULL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: NULL
     filtered: NULL
        Extra: No tables used
3 rows in set (0.00 sec)
```

```
mysql> SELECT * FROM t2;
+------+
| c1   |
+------+
|    5 |
+------+
1 row in set (0.00 sec)
```

This also means that an `EXPLAIN SELECT` statement such as the one shown here may take a long time to execute because the `BENCHMARK()` function is executed once for each row in `t1`:

```
EXPLAIN SELECT * FROM t1 AS a1, (SELECT BENCHMARK(1000000, MD5(NOW())));
```

13.2.11.9 WITH Syntax (Common Table Expressions)

A common table expression (CTE) is a named temporary result set that exists within the scope of a single statement and that can be referred to later within that statement, possibly multiple times. The following discussion describes how to write statements that use CTEs.

- Common Table Expression Syntax
- Recursive Common Table Expressions
- Recursive Common Table Expression Examples
- Common Table Expressions Compared to Similar Constructs

For information about CTE optimization, see Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

Common Table Expression Syntax

To specify common table expressions, use a `WITH` clause that has one or more comma-separated subclauses, each of which provides a subquery that produces a result set and a name associated with the subquery. The following example defines CTEs named `cte1` and `cte2` in the `WITH` clause, and refers to them in the top-level `SELECT` that follows the `WITH` clause:

```
WITH
  cte1 AS (SELECT a, b FROM table1),
  cte2 AS (SELECT c, d FROM table2)
SELECT b, d FROM cte1 JOIN cte2
WHERE cte1.a = cte2.c;
```

In the statement containing the `WITH` clause, each CTE name can be referenced to access the corresponding CTE result set.

A CTE name can be referenced in other CTEs, enabling CTEs to be defined based on other CTEs.

A CTE can refer to itself to define a recursive CTE. Common applications of recursive CTEs include series generation and traversal of hierarchical or tree-structured data.

Common table expressions are an optional part of the syntax for DML statements. They are defined using a `WITH` clause:

```
with_clause:
    WITH [RECURSIVE]
        cte_name [(col_name [, col_name] ...)] AS (subquery)
        [, cte_name [(col_name [, col_name] ...)] AS (subquery)] ...
```

`cte_name` names a single common table expression and can be used as a table reference in the statement containing the `WITH` clause.

The *subquery* part of AS (*subquery*) is called the “subquery of the CTE” and is what produces the CTE result set. The parentheses following AS are required.

A common table expression is recursive if its subquery refers to its own name. The RECURSIVE keyword must be included if any CTE in the WITH clause is recursive. For more information, see Recursive Common Table Expressions.

Determination of column names for a given CTE occurs as follows:

- If a parenthesized list of names follows the CTE name, those names are the column names:

```
WITH cte (col1, col2) AS
(
  SELECT 1, 2
  UNION ALL
  SELECT 3, 4
)
SELECT col1, col2 FROM cte;
```

  The number of names in the list must be the same as the number of columns in the result set.

- Otherwise, the column names come from the select list of the first SELECT within the AS (*subquery*) part:

```
WITH cte AS
(
  SELECT 1 AS col1, 2 AS col2
  UNION ALL
  SELECT 3, 4
)
SELECT col1, col2 FROM cte;
```

A WITH clause is permitted in these contexts:

- At the beginning of SELECT, UPDATE, and DELETE statements.

```
WITH ... SELECT ...
WITH ... UPDATE ...
WITH ... DELETE ...
```

- At the beginning of subqueries (including derived table subqueries):

```
SELECT ... WHERE id IN (WITH ... SELECT ...) ...
SELECT * FROM (WITH ... SELECT ...) AS dt ...
```

- Immediately preceding SELECT for statements that include a SELECT statement:

```
INSERT ... WITH ... SELECT ...
REPLACE ... WITH ... SELECT ...
CREATE TABLE ... WITH ... SELECT ...
CREATE VIEW ... WITH ... SELECT ...
DECLARE CURSOR ... WITH ... SELECT ...
EXPLAIN ... WITH ... SELECT ...
```

Only one WITH clause is permitted at the same level. WITH followed by WITH at the same level is not permitted, so this is illegal:

```
WITH cte1 AS (...) WITH cte2 AS (...) SELECT ...
```

However, a statement can contain multiple WITH clauses if they occur at different levels:

```
WITH cte1 AS (SELECT 1)
```

```
SELECT * FROM (WITH cte2 AS (SELECT 2) SELECT * FROM cte2 JOIN cte1) AS dt;
```

A `WITH` clause can define one or more common table expressions, but each CTE name must be unique to the clause. This is illegal:

```
WITH cte1 AS (...), cte1 AS (...) SELECT ...
```

To make the statement legal, define the CTEs with unique names:

```
WITH cte1 AS (...), cte2 AS (...) SELECT ...
```

A CTE can refer to itself or to other CTEs:

- A self-referencing CTE is recursive.
- A CTE can refer to CTEs defined earlier in the same `WITH` clause, but not those defined later.

  This constraint rules out mutually-recursive CTEs, where `cte1` references `cte2` and `cte2` references `cte1` because one of those references must be to a CTE defined later.
- A CTE in a given query block can refer to CTEs defined in query blocks at a more outer level, but not CTEs defined in query blocks at a more inner level.

For resolving references to objects with the same names, derived tables hide CTEs; and CTEs hide base tables, `TEMPORARY` tables, and views. Name resolution occurs by searching for objects in the same query block, then proceeding to outer blocks in turn while no object with the name is found.

A CTE cannot contain outer references. As with derived tables, which also prohibit outer references, this is a MySQL restriction, not a restriction of the SQL standard. For additional syntax considerations specific to recursive CTEs, see Recursive Common Table Expressions.

Recursive Common Table Expressions

A recursive common table expression is one having a subquery that refers to its own name. For example:

```
WITH RECURSIVE cte (n) AS
(
  SELECT 1
  UNION ALL
  SELECT n + 1 FROM cte WHERE n < 5
)
SELECT * FROM cte;
```

When executed, the statement produces this result, a single column containing a simple linear sequence:

```
+------+
| n    |
+------+
|    1 |
|    2 |
|    3 |
|    4 |
|    5 |
+------+
```

A recursive CTE has this structure:

- The `WITH` clause must begin with `WITH RECURSIVE` if any CTE in the `WITH` clause refers to itself. (If no CTE refers to itself, `RECURSIVE` is permitted but not required.)

  If you forget `RECURSIVE` for a recursive CTE, this error is a likely result:

```
ERROR 1146 (42S02): Table 'cte_name' doesn't exist
```

- The recursive CTE subquery has two parts, separated by `UNION [ALL]` or `UNION DISTINCT`:

```
SELECT ...      -- return initial row set
UNION ALL
SELECT ...      -- return additional row sets
```

  The first `SELECT` produces the initial row or rows for the CTE and does not refer to the CTE name. The second `SELECT` produces additional rows and recurses by referring to the CTE name in its `FROM` clause. Recursion ends when this part produces no new rows. Thus, a recursive CTE consists of a nonrecursive `SELECT` part followed by a recursive `SELECT` part.

  Either `SELECT` part can itself be a union of multiple `SELECT` statements.

- The types of the CTE result columns are inferred from the column types of the nonrecursive `SELECT` part only, and the columns are all nullable. For type determination, the recursive `SELECT` part is ignored.
- If the nonrecursive and recursive parts are separated by `UNION DISTINCT`, duplicate rows are eliminated. This is useful for queries that perform transitive closures, to avoid infinite loops.
- Each iteration of the recursive part operates only on the rows produced by the previous iteration. If the recursive part has multiple query blocks, they all operate for each iteration on the rows produced by the previous iteration, and their results are concatenated for the next iteration.

The recursive CTE subquery shown earlier has this nonrecursive part that retrieves a single row to produce the initial row set:

```
SELECT 1
```

The CTE subquery also has this recursive part:

```
SELECT n + 1 FROM cte WHERE n < 5
```

At each iteration, that `SELECT` produces a row with a new value one greater than the value of `n` from the previous row set. The first iteration operates on the initial row set (`1`) and produces `1+1=2`; the second iteration operates on the first iteration's row set (`2`) and produces `2+1=3`; and so forth. This continues until recursion ends, when `n` is no longer less than 5.

If the recursive part of a CTE produces wider values for a column than the nonrecursive part, it may be necessary to widen the column in the nonrecursive part to avoid data truncation. Consider this statement:

```
WITH RECURSIVE cte AS
(
  SELECT 1 AS n, 'abc' AS str
  UNION ALL
  SELECT n + 1, CONCAT(str, str) FROM cte WHERE n < 3
)
SELECT * FROM cte;
```

In nonstrict SQL mode, the statement produces this output:

```
+------+------+
| n    | str  |
+------+------+
|    1 | abc  |
|    2 | abc  |
```

```
|    3 | abc  |
+------+------+
```

The `str` column values are all `'abc'` because the nonrecursive `SELECT` determines the column widths. Consequently, the wider `str` values produced by the recursive `SELECT` are truncated.

In strict SQL mode, the statement produces an error:

```
ERROR 1406 (22001): Data too long for column 'str' at row 1
```

To address this issue, use `CAST()` in the nonrecursive `SELECT` to make the `str` column wider:

```
WITH RECURSIVE cte AS
(
  SELECT 1 AS n, CAST('abc' AS CHAR(20)) AS str
  UNION ALL
  SELECT n + 1, CONCAT(str, str) FROM cte WHERE n < 3
)
SELECT * FROM cte;
```

Now the statement produces this result, without truncation:

```
+------+--------------+
| n    | str          |
+------+--------------+
|    1 | abc          |
|    2 | abcabc       |
|    3 | abcabcabcabc |
+------+--------------+
```

Columns are accessed by name, not position, which means that columns in the recursive part can access columns in the nonrecursive part that have a different position, as this CTE illustrates:

```
WITH RECURSIVE cte AS
(
  SELECT 1 AS n, 1 AS p, -1 AS q
  UNION ALL
  SELECT n + 1, q * 2, p * 2 FROM cte WHERE n < 5
)
SELECT * FROM cte;
```

Because `p` in one row is derived from `q` in the previous row, and vice versa, the positive and negative values values swap positions in each successive row of the output:

```
+------+------+------+
| n    | p    | q    |
+------+------+------+
|    1 |    1 |   -1 |
|    2 |   -2 |    2 |
|    3 |    4 |   -4 |
|    4 |   -8 |    8 |
|    5 |   16 |  -16 |
+------+------+------+
```

Some syntax constraints apply within recursive CTE subqueries:

- The recursive `SELECT` part must not contain these constructs:
  - Aggregate functions such as `SUM()`
  - `GROUP BY`
  - `ORDER BY`

  - `LIMIT`
  - `DISTINCT`

  This constraint does not apply to the nonrecursive `SELECT` part of a recursive CTE. The prohibition on `DISTINCT` applies only to `UNION` members; `UNION DISTINCT` is permitted.

- The recursive `SELECT` part must reference the CTE only once and only in its `FROM` clause, not in any subquery. It can reference tables other than the CTE and join them with the CTE. If used in a join like this, the CTE must not be on the right side of a `LEFT JOIN`.

These constraints come from the SQL standard, other than the MySQL-specific exclusions of `ORDER BY`, `LIMIT`, and `DISTINCT`.

It is important for recursive CTEs that the recursive `SELECT` part include a condition to terminate recursion. Otherwise, recursion terminates only if some kind of resource exhaustion occurs. As a development technique to guard against a runaway recursive CTE, you can force termination by placing a limit on execution time:

- The `max_execution_time` system variable enforces an execution timeout for all `SELECT` statements executed within the current session.
- The `MAX_EXECUTION_TIME` optimizer hint enforces a per-query execution timeout for the `SELECT` statement in which it appears.

Suppose that a recursive CTE is mistakenly written with no recursion execution termination condition:

```
WITH RECURSIVE cte (n) AS
(
  SELECT 1
  UNION ALL
  SELECT n + 1 FROM cte
)
SELECT * FROM cte;
```

To guard against infinite recursion, set a per-session timeout by executing a statement like this prior to executing the CTE statement:

```
SET max_execution_time = 1000; -- impose one second timeout
```

Alternatively, include an optimizer hint within the CTE statement itself:

```
WITH RECURSIVE cte (n) AS
(
  SELECT 1
  UNION ALL
  SELECT n + 1 FROM cte
)
SELECT /*+ MAX_EXECUTION_TIME(1000) */ * FROM cte;
```

If a recursive query without an execution time limit enters an infinite loop, you can terminate it from another session using `KILL QUERY`. Within the session itself, the client program used to run the query might provide a way to kill the query. For example, in `mysql`, typing **Control+C** interrupts the current statement.

For recursive CTEs, `EXPLAIN` output rows for recursive `SELECT` parts display `Recursive` in the `Extra` column.

Cost estimates displayed by `EXPLAIN` represent cost per iteration, which might differ considerably from total cost. The optimizer cannot predict the number of iterations because it cannot predict when the `WHERE` clause will become false.

CTE actual cost may also be affected by result set size. A CTE that produces many rows may require an internal temporary table large enough to be converted from in-memory to on-disk format and may suffer a performance penalty. If so, increasing the permitted in-memory temporary table size may improve performance; see Section 8.4.4, “Internal Temporary Table Use in MySQL”.

Recursive Common Table Expression Examples

As mentioned previously, recursive common table expressions (CTEs) are frequently used for series generation and traversing hierarchical or tree-structured data. This section shows some simple examples of these techniques.

- Fibonacci Series Generation
- Date Series Generation
- Hierarchical Data Traversal

Fibonacci Series Generation

A Fibonacci series begins with the two numbers 0 and 1 (or 1 and 1) and each number after that is the sum of the previous two numbers. A recursive common table expression can generate a Fibonacci series if each row produced by the recursive `SELECT` has access to the two previous numbers from the series. The following CTE generates a 10-number series using 0 and 1 as the first two numbers:

```
WITH RECURSIVE fibonacci (n, fib_n, next_fib_n) AS
(
  SELECT 1, 0, 1
  UNION ALL
  SELECT n + 1, next_fib_n, fib_n + next_fib_n
    FROM fibonacci WHERE n < 10
)
SELECT * FROM fibonacci;
```

The CTE produces this result:

```
+------+-------+------------+
| n    | fib_n | next_fib_n |
+------+-------+------------+
|    1 |     0 |          1 |
|    2 |     1 |          1 |
|    3 |     1 |          2 |
|    4 |     2 |          3 |
|    5 |     3 |          5 |
|    6 |     5 |          8 |
|    7 |     8 |         13 |
|    8 |    13 |         21 |
|    9 |    21 |         34 |
|   10 |    34 |         55 |
+------+-------+------------+
```

How the CTE works:

- `n` is a display column to indicate that the row contains the `n`-th Fibonacci number. For example, the 8th Fibonacci number is 13.
- The `fib_n` column displays Fibonacci number `n`.
- The `next_fib_n` column displays the next Fibonacci number after number `n`. This column provides the next series value to the next row, so that row can produce the sum of the two previous series values in its `fib_n` column.
- Recursion ends when `n` reaches 10. This is an arbitrary choice, to limit the output to a small set of rows.

The preceding output shows the entire CTE result. To select just part of it, add an appropriate `WHERE` clause to the top-level `SELECT`. For example, to select the 8th Fibonacci number, do this:

```
mysql> WITH RECURSIVE fibonacci ...
       ...
       SELECT fib_n FROM fibonacci WHERE n = 8;
+-------+
| fib_n |
+-------+
|    13 |
+-------+
```

Date Series Generation

A common table expression can generate a series of successive dates, which is useful for generating summaries that include a row for all dates in the series, including dates not represented in the summarized data.

Suppose that a table of sales numbers contains these rows:

```
mysql> SELECT * FROM sales ORDER BY date, price;
+------------+--------+
| date       | price  |
+------------+--------+
| 2017-01-03 | 100.00 |
| 2017-01-03 | 200.00 |
| 2017-01-06 |  50.00 |
| 2017-01-08 |  10.00 |
| 2017-01-08 |  20.00 |
| 2017-01-08 | 150.00 |
| 2017-01-10 |   5.00 |
+------------+--------+
```

This query summarizes the sales per day:

```
mysql> SELECT date, SUM(price) AS sum_price
       FROM sales
       GROUP BY date
       ORDER BY date;
+------------+-----------+
| date       | sum_price |
+------------+-----------+
| 2017-01-03 |    300.00 |
| 2017-01-06 |     50.00 |
| 2017-01-08 |    180.00 |
| 2017-01-10 |      5.00 |
+------------+-----------+
```

However, that result contains “holes” for dates not represented in the range of dates spanned by the table. A result that represents all dates in the range can be produces using a recursive CTE to generate that set of dates, joined with a `LEFT JOIN` to the sales data.

Here is the CTE to generate the date range series:

```
WITH RECURSIVE dates (date) AS
(
  SELECT MIN(date) FROM sales
  UNION ALL
  SELECT date + INTERVAL 1 DAY FROM dates
  WHERE date + INTERVAL 1 DAY <= (SELECT MAX(date) FROM sales)
)
SELECT * FROM dates;
```

The CTE produces this result:

```
+------------+
| date       |
+------------+
| 2017-01-03 |
| 2017-01-04 |
| 2017-01-05 |
| 2017-01-06 |
| 2017-01-07 |
| 2017-01-08 |
| 2017-01-09 |
| 2017-01-10 |
+------------+
```

How the CTE works:

- The nonrecursive `SELECT` produces the lowest date in the date range spanned by the `sales` table.
- Each row produced by the recursive `SELECT` adds one day to the date produced by the previous row.
- Recursion ends after the dates reach the highest date in the date range spanned by the `sales` table.

Joining the CTE with a `LEFT JOIN` against the `sales` table produces the sales summary with a row for each date in the range:

```
WITH RECURSIVE dates (date) AS
(
  SELECT MIN(date) FROM sales
  UNION ALL
  SELECT date + INTERVAL 1 DAY FROM dates
  WHERE date + INTERVAL 1 DAY <= (SELECT MAX(date) FROM sales)
)
SELECT dates.date, COALESCE(SUM(price), 0) AS sum_price
FROM dates LEFT JOIN sales ON dates.date = sales.date
GROUP BY dates.date
ORDER BY dates.date;
```

The output looks like this:

```
+------------+-----------+
| date       | sum_price |
+------------+-----------+
| 2017-01-03 |    300.00 |
| 2017-01-04 |      0.00 |
| 2017-01-05 |      0.00 |
| 2017-01-06 |     50.00 |
| 2017-01-07 |      0.00 |
| 2017-01-08 |    180.00 |
| 2017-01-09 |      0.00 |
| 2017-01-10 |      5.00 |
+------------+-----------+
```

Some points to note:

- Are the queries inefficient, particularly the one with the `MAX()` subquery executed for each row in the recursive `SELECT`? Checking with `EXPLAIN` shows that the subqueries are optimized away for efficiency.
- The use of `COALESCE()` avoids displaying `NULL` in the `sum_price` column on days for which no sales data occur in the `sales` table.

Hierarchical Data Traversal

Recursive common table expressions are useful for traversing data that forms a hierarchy. Consider these statements that create a small data set that shows, for each employee in a company, the

employee name and ID number, and the ID of the employee's manager. The top-level employee (the CEO), has a manager ID of `NULL` (no manager).

```
CREATE TABLE employees (
  id         INT PRIMARY KEY NOT NULL,
  name       VARCHAR(100) NOT NULL,
  manager_id INT NULL,
  INDEX (manager_id),
FOREIGN KEY (manager_id) REFERENCES EMPLOYEES (id)
);
INSERT INTO employees VALUES
(333, "Yasmina", NULL),  # Yasmina is the CEO (manager_id is NULL)
(198, "John", 333),      # John has ID 198 and reports to 333 (Yasmina)
(692, "Tarek", 333),
(29, "Pedro", 198),
(4610, "Sarah", 29),
(72, "Pierre", 29),
(123, "Adil", 692);
```

The resulting data set looks like this:

```
mysql> SELECT * FROM employees ORDER BY id;
+------+---------+------------+
| id   | name    | manager_id |
+------+---------+------------+
|   29 | Pedro   |        198 |
|   72 | Pierre  |         29 |
|  123 | Adil    |        692 |
|  198 | John    |        333 |
|  333 | Yasmina |       NULL |
|  692 | Tarek   |        333 |
| 4610 | Sarah   |         29 |
+------+---------+------------+
```

To produce the organizational chart with the management chain for each employee (that is, the path from CEO to employee), use a recursive CTE:

```
WITH RECURSIVE employee_paths (id, name, path) AS
(
  SELECT id, name, CAST(id AS CHAR(200))
    FROM employees
    WHERE manager_id IS NULL
  UNION ALL
  SELECT e.id, e.name, CONCAT(ep.path, ',', e.id)
    FROM employee_paths AS ep JOIN employees AS e
      ON ep.id = e.manager_id
)
SELECT * FROM employee_paths ORDER BY path;
```

The CTE produces this output:

```
+------+---------+-----------------+
| id   | name    | path            |
+------+---------+-----------------+
|  333 | Yasmina | 333             |
|  198 | John    | 333,198         |
|   29 | Pedro   | 333,198,29      |
| 4610 | Sarah   | 333,198,29,4610 |
|   72 | Pierre  | 333,198,29,72   |
|  692 | Tarek   | 333,692         |
|  123 | Adil    | 333,692,123     |
+------+---------+-----------------+
```

How the CTE works:

- The nonrecursive `SELECT` produces the row for the CEO (the row with a `NULL` manager ID).

The `path` column is widened to `CHAR(200)` to ensure that there is room for the longer `path` values produced by the recursive `SELECT`.

- Each row produced by the recursive `SELECT` finds all employees who report directly to an employee produced by a previous row. For each such employee, the row includes the employee ID and name, and the employee management chain. The chain is the manager's chain, with the employee ID added to the end.
- Recursion ends when employees have no others who report to them.

To find the path for a specific employee or employees, add a `WHERE` clause to the top-level `SELECT`. For example, to display the results for Tarek and Sarah, modify that `SELECT` like this:

```
mysql> WITH RECURSIVE ...
       ...
       SELECT * FROM employees_extended
       WHERE id IN (692, 4610)
       ORDER BY path;
+------+-------+-----------------+
| id   | name  | path            |
+------+-------+-----------------+
| 4610 | Sarah | 333,198,29,4610 |
|  692 | Tarek | 333,692         |
+------+-------+-----------------+
```

Common Table Expressions Compared to Similar Constructs

Common table expressions (CTEs) are similar to derived tables in some ways:

- Both constructs are named.
- Both constructs exist for the scope of a single statement.

Because of these similarities, CTEs and derived tables often can be used interchangeably. As a trivial example, these statements are equivalent:

```
WITH cte AS (SELECT 1) SELECT * FROM cte;
SELECT * FROM (SELECT 1) AS dt;
```

However, CTEs have some advantages over derived tables:

- A derived table can be referenced only a single time within a query. A CTE can be referenced multiple times. To use multiple instances of a derived table result, you must derive the result multiple times.
- A CTE can be self-referencing (recursive).
- One CTE can refer to another.
- A CTE may be easier to read when its definition appears at the beginning of the statement rather than embedded within it.

CTEs are similar to tables created with `CREATE [TEMPORARY] TABLE` but need not be defined or dropped explicitly. For a CTE, you need no privileges to create tables.

13.2.11.10 Subquery Errors

There are some errors that apply only to subqueries. This section describes them.

- Unsupported subquery syntax:

```
ERROR 1235 (ER_NOT_SUPPORTED_YET)
SQLSTATE = 42000
Message = "This version of MySQL doesn't yet support
'LIMIT & IN/ALL/ANY/SOME subquery'"
```

This means that MySQL does not support statements of the following form:

```
SELECT * FROM t1 WHERE s1 IN (SELECT s2 FROM t2 ORDER BY s1 LIMIT 1)
```

- Incorrect number of columns from subquery:

```
ERROR 1241 (ER_OPERAND_COL)
SQLSTATE = 21000
Message = "Operand should contain 1 column(s)"
```

This error occurs in cases like this:

```
SELECT (SELECT column1, column2 FROM t2) FROM t1;
```

You may use a subquery that returns multiple columns, if the purpose is row comparison. In other contexts, the subquery must be a scalar operand. See Section 13.2.11.5, “Row Subqueries”.

- Incorrect number of rows from subquery:

```
ERROR 1242 (ER_SUBSELECT_NO_1_ROW)
SQLSTATE = 21000
Message = "Subquery returns more than 1 row"
```

This error occurs for statements where the subquery must return at most one row but returns multiple rows. Consider the following example:

```
SELECT * FROM t1 WHERE column1 = (SELECT column1 FROM t2);
```

If `SELECT column1 FROM t2` returns just one row, the previous query will work. If the subquery returns more than one row, error 1242 will occur. In that case, the query should be rewritten as:

```
SELECT * FROM t1 WHERE column1 = ANY (SELECT column1 FROM t2);
```

- Incorrectly used table in subquery:

```
Error 1093 (ER_UPDATE_TABLE_USED)
SQLSTATE = HY000
Message = "You can't specify target table 'x'
for update in FROM clause"
```

This error occurs in cases such as the following, which attempts to modify a table and select from the same table in the subquery:

```
UPDATE t1 SET column2 = (SELECT MAX(column1) FROM t1);
```

You can use a subquery for assignment within an `UPDATE` statement because subqueries are legal in `UPDATE` and `DELETE` statements as well as in `SELECT` statements. However, you cannot use the same table (in this case, table `t1`) for both the subquery `FROM` clause and the update target.

For transactional storage engines, the failure of a subquery causes the entire statement to fail. For nontransactional storage engines, data modifications made before the error was encountered are preserved.

13.2.11.11 Optimizing Subqueries

Development is ongoing, so no optimization tip is reliable for the long term. The following list provides some interesting tricks that you might want to play with. See also Section 8.2.2, “Optimizing Subqueries, Derived Tables, View References, and Common Table Expressions”.

- Use subquery clauses that affect the number or order of the rows in the subquery. For example:

```
SELECT * FROM t1 WHERE t1.column1 IN
  (SELECT column1 FROM t2 ORDER BY column1);
SELECT * FROM t1 WHERE t1.column1 IN
  (SELECT DISTINCT column1 FROM t2);
SELECT * FROM t1 WHERE EXISTS
  (SELECT * FROM t2 LIMIT 1);
```

- Replace a join with a subquery. For example, try this:

```
SELECT DISTINCT column1 FROM t1 WHERE t1.column1 IN (
  SELECT column1 FROM t2);
```

  Instead of this:

```
SELECT DISTINCT t1.column1 FROM t1, t2
  WHERE t1.column1 = t2.column1;
```

- Some subqueries can be transformed to joins for compatibility with older versions of MySQL that do not support subqueries. However, in some cases, converting a subquery to a join may improve performance. See Section 13.2.11.12, “Rewriting Subqueries as Joins”.

- Move clauses from outside to inside the subquery. For example, use this query:

```
SELECT * FROM t1
  WHERE s1 IN (SELECT s1 FROM t1 UNION ALL SELECT s1 FROM t2);
```

  Instead of this query:

```
SELECT * FROM t1
  WHERE s1 IN (SELECT s1 FROM t1) OR s1 IN (SELECT s1 FROM t2);
```

  For another example, use this query:

```
SELECT (SELECT column1 + 5 FROM t1) FROM t2;
```

  Instead of this query:

```
SELECT (SELECT column1 FROM t1) + 5 FROM t2;
```

- Use a row subquery instead of a correlated subquery. For example, use this query:

```
SELECT * FROM t1
  WHERE (column1,column2) IN (SELECT column1,column2 FROM t2);
```

  Instead of this query:

```
SELECT * FROM t1
  WHERE EXISTS (SELECT * FROM t2 WHERE t2.column1=t1.column1
                AND t2.column2=t1.column2);
```

- Use `NOT (a = ANY (...))` rather than `a <> ALL (...)`.

- Use `x = ANY (table containing (1,2))` rather than `x=1 OR x=2`.
- Use `= ANY` rather than `EXISTS`.
- For uncorrelated subqueries that always return one row, `IN` is always slower than `=`. For example, use this query:

```
SELECT * FROM t1
  WHERE t1.col_name = (SELECT a FROM t2 WHERE b = some_const);
```

  Instead of this query:

```
SELECT * FROM t1
  WHERE t1.col_name IN (SELECT a FROM t2 WHERE b = some_const);
```

These tricks might cause programs to go faster or slower. Using MySQL facilities like the `BENCHMARK()` function, you can get an idea about what helps in your own situation. See Section 12.14, “Information Functions”.

Some optimizations that MySQL itself makes are:

- MySQL executes uncorrelated subqueries only once. Use `EXPLAIN` to make sure that a given subquery really is uncorrelated.
- MySQL rewrites `IN`, `ALL`, `ANY`, and `SOME` subqueries in an attempt to take advantage of the possibility that the select-list columns in the subquery are indexed.
- MySQL replaces subqueries of the following form with an index-lookup function, which `EXPLAIN` describes as a special join type (`unique_subquery` or `index_subquery`):

```
... IN (SELECT indexed_column FROM single_table ...)
```

- MySQL enhances expressions of the following form with an expression involving `MIN()` or `MAX()`, unless `NULL` values or empty sets are involved:

```
value {ALL|ANY|SOME} {> | < | >= | <=} (uncorrelated subquery)
```

  For example, this `WHERE` clause:

```
WHERE 5 > ALL (SELECT x FROM t)
```

  might be treated by the optimizer like this:

```
WHERE 5 > (SELECT MAX(x) FROM t)
```

See also MySQL Internals: How MySQL Transforms Subqueries.

13.2.11.12 Rewriting Subqueries as Joins

Sometimes there are other ways to test membership in a set of values than by using a subquery. Also, on some occasions, it is not only possible to rewrite a query without a subquery, but it can be more efficient to make use of some of these techniques rather than to use subqueries. One of these is the `IN()` construct:

For example, this query:

```
SELECT * FROM t1 WHERE id IN (SELECT id FROM t2);
```

Can be rewritten as:

```
SELECT DISTINCT t1.* FROM t1, t2 WHERE t1.id=t2.id;
```

The queries:

```
SELECT * FROM t1 WHERE id NOT IN (SELECT id FROM t2);
SELECT * FROM t1 WHERE NOT EXISTS (SELECT id FROM t2 WHERE t1.id=t2.id);
```

Can be rewritten as:

```
SELECT table1.*
  FROM table1 LEFT JOIN table2 ON table1.id=table2.id
  WHERE table2.id IS NULL;
```

A `LEFT [OUTER] JOIN` can be faster than an equivalent subquery because the server might be able to optimize it better—a fact that is not specific to MySQL Server alone. Prior to SQL-92, outer joins did not exist, so subqueries were the only way to do certain things. Today, MySQL Server and many other modern database systems offer a wide range of outer join types.

MySQL Server supports multiple-table `DELETE` statements that can be used to efficiently delete rows based on information from one table or even from many tables at the same time. Multiple-table `UPDATE` statements are also supported. See Section 13.2.2, “DELETE Syntax”, and Section 13.2.12, “UPDATE Syntax”.

13.2.12 UPDATE Syntax

`UPDATE` is a DML statement that modifies rows in a table.

An `UPDATE` statement can start with a `WITH` clause to define common table expressions accessible within the `UPDATE`. See Section 13.2.11.9, “WITH Syntax (Common Table Expressions)”.

Single-table syntax:

```
UPDATE [LOW_PRIORITY] [IGNORE] table_reference
    SET col_name1={expr1|DEFAULT} [, col_name2={expr2|DEFAULT}] ...
    [WHERE where_condition]
    [ORDER BY ...]
    [LIMIT row_count]
```

Multiple-table syntax:

```
UPDATE [LOW_PRIORITY] [IGNORE] table_references
    SET col_name1={expr1|DEFAULT} [, col_name2={expr2|DEFAULT}] ...
    [WHERE where_condition]
```

For the single-table syntax, the `UPDATE` statement updates columns of existing rows in the named table with new values. The `SET` clause indicates which columns to modify and the values they should be given. Each value can be given as an expression, or the keyword `DEFAULT` to set a column explicitly to its default value. The `WHERE` clause, if given, specifies the conditions that identify which rows to update. With no `WHERE` clause, all rows are updated. If the `ORDER BY` clause is specified, the rows are updated in the order that is specified. The `LIMIT` clause places a limit on the number of rows that can be updated.

For the multiple-table syntax, `UPDATE` updates rows in each table named in *`table_references`* that satisfy the conditions. Each matching row is updated once, even if it matches the conditions multiple times. For multiple-table syntax, `ORDER BY` and `LIMIT` cannot be used.

For partitioned tables, both the single-single and multiple-table forms of this statement support the use of a `PARTITION` option as part of a table reference. This option takes a list of one or more partitions

or subpartitions (or both). Only the partitions (or subpartitions) listed are checked for matches, and a row that is not in any of these partitions or subpartitions is not updated, whether it satisfies the *`where_condition`* or not.

Note

Unlike the case when using `PARTITION` with an `INSERT` or `REPLACE` statement, an otherwise valid `UPDATE ... PARTITION` statement is considered successful even if no rows in the listed partitions (or subpartitions) match the *`where_condition`*.

See Section 22.5, “Partition Selection”, for more information and examples.

`where_condition` is an expression that evaluates to true for each row to be updated. For expression syntax, see Section 9.5, “Expression Syntax”.

`table_references` and *`where_condition`* are specified as described in Section 13.2.10, “SELECT Syntax”.

You need the `UPDATE` privilege only for columns referenced in an `UPDATE` that are actually updated. You need only the `SELECT` privilege for any columns that are read but not modified.

The `UPDATE` statement supports the following modifiers:

- With the `LOW_PRIORITY` modifier, execution of the `UPDATE` is delayed until no other clients are reading from the table. This affects only storage engines that use only table-level locking (such as `MyISAM`, `MEMORY`, and `MERGE`).
- With the `IGNORE` modifier, the update statement does not abort even if errors occur during the update. Rows for which duplicate-key conflicts occur on a unique key value are not updated. Rows updated to values that would cause data conversion errors are updated to the closest valid values instead. For more information, see Comparison of the IGNORE Keyword and Strict SQL Mode.

`UPDATE IGNORE` statements, including those having an `ORDER BY` clause, are flagged as unsafe for statement-based replication. (This is because the order in which the rows are updated determines which rows are ignored.) With this change, such statements produce a warning in the log when using statement-based mode and are logged using the row-based format when using `MIXED` mode. (Bug #11758262, Bug #50439) See Section 18.2.1.3, “Determination of Safe and Unsafe Statements in Binary Logging”, for more information.

If you access a column from the table to be updated in an expression, `UPDATE` uses the current value of the column. For example, the following statement sets `col1` to one more than its current value:

```
UPDATE t1 SET col1 = col1 + 1;
```

The second assignment in the following statement sets `col2` to the current (updated) `col1` value, not the original `col1` value. The result is that `col1` and `col2` have the same value. This behavior differs from standard SQL.

```
UPDATE t1 SET col1 = col1 + 1, col2 = col1;
```

Single-table `UPDATE` assignments are generally evaluated from left to right. For multiple-table updates, there is no guarantee that assignments are carried out in any particular order.

If you set a column to the value it currently has, MySQL notices this and does not update it.

If you update a column that has been declared `NOT NULL` by setting to `NULL`, an error occurs if strict SQL mode is enabled; otherwise, the column is set to the implicit default value for the column data type and the warning count is incremented. The implicit default value is `0` for numeric types, the empty string

('') for string types, and the "zero" value for date and time types. See Section 11.7, "Data Type Default Values".

If a generated column is updated explicitly, the only permitted value is DEFAULT. For information about generated columns, see Section 13.1.15.8, "CREATE TABLE and Generated Columns".

UPDATE returns the number of rows that were actually changed. The mysql_info() C API function returns the number of rows that were matched and updated and the number of warnings that occurred during the UPDATE.

You can use LIMIT *row_count* to restrict the scope of the UPDATE. A LIMIT clause is a rows-matched restriction. The statement stops as soon as it has found *row_count* rows that satisfy the WHERE clause, whether or not they actually were changed.

If an UPDATE statement includes an ORDER BY clause, the rows are updated in the order specified by the clause. This can be useful in certain situations that might otherwise result in an error. Suppose that a table t contains a column id that has a unique index. The following statement could fail with a duplicate-key error, depending on the order in which rows are updated:

```
UPDATE t SET id = id + 1;
```

For example, if the table contains 1 and 2 in the id column and 1 is updated to 2 before 2 is updated to 3, an error occurs. To avoid this problem, add an ORDER BY clause to cause the rows with larger id values to be updated before those with smaller values:

```
UPDATE t SET id = id + 1 ORDER BY id DESC;
```

You can also perform UPDATE operations covering multiple tables. However, you cannot use ORDER BY or LIMIT with a multiple-table UPDATE. The *table_references* clause lists the tables involved in the join. Its syntax is described in Section 13.2.10.2, "JOIN Syntax". Here is an example:

```
UPDATE items,month SET items.price=month.price
WHERE items.id=month.id;
```

The preceding example shows an inner join that uses the comma operator, but multiple-table UPDATE statements can use any type of join permitted in SELECT statements, such as LEFT JOIN.

If you use a multiple-table UPDATE statement involving InnoDB tables for which there are foreign key constraints, the MySQL optimizer might process tables in an order that differs from that of their parent/child relationship. In this case, the statement fails and rolls back. Instead, update a single table and rely on the ON UPDATE capabilities that InnoDB provides to cause the other tables to be modified accordingly. See Section 15.8.1.6, "InnoDB and FOREIGN KEY Constraints".

You cannot update a table and select from the same table in a subquery.

In MySQL 8.0, an UPDATE on a partitioned table using a storage engine such as MyISAM that employs table-level locks locks only those partitions containing rows that match the UPDATE statement's WHERE clause, as long as none of the table's partitioning columns are updated. (For storage engines such as InnoDB that employ row-level locking, no locking of partitions takes place.) For more information, see Partitioning and Locking.

13.3 Transactional and Locking Statements

MySQL supports local transactions (within a given client session) through statements such as SET autocommit, START TRANSACTION, COMMIT, and ROLLBACK. See Section 13.3.1, "START TRANSACTION, COMMIT, and ROLLBACK Syntax". XA transaction support enables MySQL to participate in distributed transactions as well. See Section 13.3.7, "XA Transactions".

13.3.1 START TRANSACTION, COMMIT, and ROLLBACK Syntax

```
START TRANSACTION
    [transaction_characteristic [, transaction_characteristic] ...]

transaction_characteristic:
    WITH CONSISTENT SNAPSHOT
  | READ WRITE
  | READ ONLY

BEGIN [WORK]
COMMIT [WORK] [AND [NO] CHAIN] [[NO] RELEASE]
ROLLBACK [WORK] [AND [NO] CHAIN] [[NO] RELEASE]
SET autocommit = {0 | 1}
```

These statements provide control over use of transactions:

- `START TRANSACTION` or `BEGIN` start a new transaction.
- `COMMIT` commits the current transaction, making its changes permanent.
- `ROLLBACK` rolls back the current transaction, canceling its changes.
- `SET autocommit` disables or enables the default autocommit mode for the current session.

By default, MySQL runs with autocommit mode enabled. This means that as soon as you execute a statement that updates (modifies) a table, MySQL stores the update on disk to make it permanent. The change cannot be rolled back.

To disable autocommit mode implicitly for a single series of statements, use the `START TRANSACTION` statement:

```
START TRANSACTION;
SELECT @A:=SUM(salary) FROM table1 WHERE type=1;
UPDATE table2 SET summary=@A WHERE type=1;
COMMIT;
```

With `START TRANSACTION`, autocommit remains disabled until you end the transaction with `COMMIT` or `ROLLBACK`. The autocommit mode then reverts to its previous state.

`START TRANSACTION` permits several modifiers that control transaction characteristics. To specify multiple modifiers, separate them by commas.

- The `WITH CONSISTENT SNAPSHOT` modifier starts a consistent read for storage engines that are capable of it. This applies only to `InnoDB`. The effect is the same as issuing a `START TRANSACTION` followed by a `SELECT` from any `InnoDB` table. See Section 15.5.2.3, “Consistent Nonlocking Reads”. The `WITH CONSISTENT SNAPSHOT` modifier does not change the current transaction isolation level, so it provides a consistent snapshot only if the current isolation level is one that permits a consistent read. The only isolation level that permits a consistent read is `REPEATABLE READ`. For all other isolation levels, the `WITH CONSISTENT SNAPSHOT` clause is ignored. A warning is generated when the `WITH CONSISTENT SNAPSHOT` clause is ignored.
- The `READ WRITE` and `READ ONLY` modifiers set the transaction access mode. They permit or prohibit changes to tables used in the transaction. The `READ ONLY` restriction prevents the transaction from modifying or locking both transactional and nontransactional tables that are visible to other transactions; the transaction can still modify or lock temporary tables.

  MySQL enables extra optimizations for queries on `InnoDB` tables when the transaction is known to be read-only. Specifying `READ ONLY` ensures these optimizations are applied in cases where the read-only status cannot be determined automatically. See Section 8.5.3, “Optimizing InnoDB Read-Only Transactions” for more information.

  If no access mode is specified, the default mode applies. Unless the default has been changed, it is read/write. It is not permitted to specify both `READ WRITE` and `READ ONLY` in the same statement.

In read-only mode, it remains possible to change tables created with the `TEMPORARY` keyword using DML statements. Changes made with DDL statements are not permitted, just as with permanent tables.

For additional information about transaction access mode, including ways to change the default mode, see Section 13.3.6, "SET TRANSACTION Syntax".

If the `read_only` system variable is enabled, explicitly starting a transaction with `START TRANSACTION READ WRITE` requires the `SUPER` privilege.

Important

Many APIs used for writing MySQL client applications (such as JDBC) provide their own methods for starting transactions that can (and sometimes should) be used instead of sending a `START TRANSACTION` statement from the client. See Chapter 27, *Connectors and APIs*, or the documentation for your API, for more information.

To disable autocommit mode explicitly, use the following statement:

```
SET autocommit=0;
```

After disabling autocommit mode by setting the `autocommit` variable to zero, changes to transaction-safe tables (such as those for `InnoDB` or `NDB`) are not made permanent immediately. You must use `COMMIT` to store your changes to disk or `ROLLBACK` to ignore the changes.

`autocommit` is a session variable and must be set for each session. To disable autocommit mode for each new connection, see the description of the `autocommit` system variable at Section 5.1.5, "Server System Variables".

`BEGIN` and `BEGIN WORK` are supported as aliases of `START TRANSACTION` for initiating a transaction. `START TRANSACTION` is standard SQL syntax, is the recommended way to start an ad-hoc transaction, and permits modifiers that `BEGIN` does not.

The `BEGIN` statement differs from the use of the `BEGIN` keyword that starts a `BEGIN ... END` compound statement. The latter does not begin a transaction. See Section 13.6.1, "BEGIN ... END Compound-Statement Syntax".

Note

Within all stored programs (stored procedures and functions, triggers, and events), the parser treats `BEGIN [WORK]` as the beginning of a `BEGIN ... END` block. Begin a transaction in this context with `START TRANSACTION` instead.

The optional `WORK` keyword is supported for `COMMIT` and `ROLLBACK`, as are the `CHAIN` and `RELEASE` clauses. `CHAIN` and `RELEASE` can be used for additional control over transaction completion. The value of the `completion_type` system variable determines the default completion behavior. See Section 5.1.5, "Server System Variables".

The `AND CHAIN` clause causes a new transaction to begin as soon as the current one ends, and the new transaction has the same isolation level as the just-terminated transaction. The `RELEASE` clause causes the server to disconnect the current client session after terminating the current transaction. Including the `NO` keyword suppresses `CHAIN` or `RELEASE` completion, which can be useful if the `completion_type` system variable is set to cause chaining or release completion by default.

Beginning a transaction causes any pending transaction to be committed. See Section 13.3.3, "Statements That Cause an Implicit Commit", for more information.

Beginning a transaction also causes table locks acquired with `LOCK TABLES` to be released, as though you had executed `UNLOCK TABLES`. Beginning a transaction does not release a global read lock acquired with `FLUSH TABLES WITH READ LOCK`.

For best results, transactions should be performed using only tables managed by a single transaction-safe storage engine. Otherwise, the following problems can occur:

- If you use tables from more than one transaction-safe storage engine (such as `InnoDB`), and the transaction isolation level is not `SERIALIZABLE`, it is possible that when one transaction commits, another ongoing transaction that uses the same tables will see only some of the changes made by the first transaction. That is, the atomicity of transactions is not guaranteed with mixed engines and inconsistencies can result. (If mixed-engine transactions are infrequent, you can use `SET TRANSACTION ISOLATION LEVEL` to set the isolation level to `SERIALIZABLE` on a per-transaction basis as necessary.)
- If you use tables that are not transaction-safe within a transaction, changes to those tables are stored at once, regardless of the status of autocommit mode.
- If you issue a `ROLLBACK` statement after updating a nontransactional table within a transaction, an `ER_WARNING_NOT_COMPLETE_ROLLBACK` warning occurs. Changes to transaction-safe tables are rolled back, but not changes to nontransaction-safe tables.

Each transaction is stored in the binary log in one chunk, upon `COMMIT`. Transactions that are rolled back are not logged. (**Exception**: Modifications to nontransactional tables cannot be rolled back. If a transaction that is rolled back includes modifications to nontransactional tables, the entire transaction is logged with a `ROLLBACK` statement at the end to ensure that modifications to the nontransactional tables are replicated.) See Section 5.4.4, “The Binary Log”.

You can change the isolation level or access mode for transactions with the `SET TRANSACTION` statement. See Section 13.3.6, “SET TRANSACTION Syntax”.

Rolling back can be a slow operation that may occur implicitly without the user having explicitly asked for it (for example, when an error occurs). Because of this, `SHOW PROCESSLIST` displays `Rolling back` in the `State` column for the session, not only for explicit rollbacks performed with the `ROLLBACK` statement but also for implicit rollbacks.

Note

In MySQL 8.0, `BEGIN`, `COMMIT`, and `ROLLBACK` are not affected by `--replicate-do-db` or `--replicate-ignore-db` rules.

13.3.2 Statements That Cannot Be Rolled Back

Some statements cannot be rolled back. In general, these include data definition language (DDL) statements, such as those that create or drop databases, those that create, drop, or alter tables or stored routines.

You should design your transactions not to include such statements. If you issue a statement early in a transaction that cannot be rolled back, and then another statement later fails, the full effect of the transaction cannot be rolled back in such cases by issuing a `ROLLBACK` statement.

13.3.3 Statements That Cause an Implicit Commit

The statements listed in this section (and any synonyms for them) implicitly end any transaction active in the current session, as if you had done a `COMMIT` before executing the statement.

Most of these statements also cause an implicit commit after executing. The intent is to handle each such statement in its own special transaction because it cannot be rolled back anyway. Transaction-control and locking statements are exceptions: If an implicit commit occurs before execution, another does not occur after.

- **Data definition language (DDL) statements that define or modify database objects.** `ALTER EVENT`, `ALTER FUNCTION`, `ALTER PROCEDURE`, `ALTER SERVER`, `ALTER TABLE`, `ALTER VIEW`, `CREATE DATABASE`, `CREATE EVENT`, `CREATE FUNCTION`, `CREATE INDEX`, `CREATE PROCEDURE`, `CREATE ROLE`, `CREATE SERVER`, `CREATE TABLE`, `CREATE TRIGGER`, `CREATE VIEW`, `DROP DATABASE`, `DROP EVENT`, `DROP FUNCTION`, `DROP INDEX`, `DROP PROCEDURE`, `DROP ROLE`, `DROP SERVER`, `DROP TABLE`, `DROP TRIGGER`, `DROP VIEW`, `INSTALL PLUGIN`, `RENAME TABLE`, `TRUNCATE TABLE`, `UNINSTALL PLUGIN`.

  `CREATE TABLE` and `DROP TABLE` statements do not commit a transaction if the `TEMPORARY` keyword is used. (This does not apply to other operations on temporary tables such as `ALTER TABLE` and `CREATE INDEX`, which do cause a commit.) However, although no implicit commit occurs, neither can the statement be rolled back, which means that the use of such statements causes transactional atomicity to be violated. For example, if you use `CREATE TEMPORARY TABLE` and then roll back the transaction, the table remains in existence.

  The `CREATE TABLE` statement in `InnoDB` is processed as a single transaction. This means that a `ROLLBACK` from the user does not undo `CREATE TABLE` statements the user made during that transaction.

  `CREATE TABLE ... SELECT` causes an implicit commit before and after the statement is executed when you are creating nontemporary tables. (No commit occurs for `CREATE TEMPORARY TABLE ... SELECT`.) This is to prevent an issue during replication where the table could be created on the master after a rollback, but fail to be recorded in the binary log, and therefore not replicated to the slave.

- **Statements that implicitly use or modify tables in the `mysql` database.** `ALTER USER`, `CREATE USER`, `DROP USER`, `GRANT`, `RENAME USER`, `REVOKE`, `SET PASSWORD`.

- **Transaction-control and locking statements.** `BEGIN`, `LOCK TABLES`, `SET autocommit = 1` (if the value is not already 1), `START TRANSACTION`, `UNLOCK TABLES`.

  `UNLOCK TABLES` commits a transaction only if any tables currently have been locked with `LOCK TABLES` to acquire nontransactional table locks. A commit does not occur for `UNLOCK TABLES` following `FLUSH TABLES WITH READ LOCK` because the latter statement does not acquire table-level locks.

  Transactions cannot be nested. This is a consequence of the implicit commit performed for any current transaction when you issue a `START TRANSACTION` statement or one of its synonyms.

  Statements that cause an implicit commit cannot be used in an XA transaction while the transaction is in an `ACTIVE` state.

  The `BEGIN` statement differs from the use of the `BEGIN` keyword that starts a `BEGIN ... END` compound statement. The latter does not cause an implicit commit. See Section 13.6.1, “BEGIN ... END Compound-Statement Syntax”.

- **Data loading statements.** `LOAD DATA INFILE`. `LOAD DATA INFILE` causes an implicit commit only for tables using the `NDB` storage engine.

- **Administrative statements.** `ANALYZE TABLE`, `CACHE INDEX`, `CHECK TABLE`, `FLUSH`, `LOAD INDEX INTO CACHE`, `OPTIMIZE TABLE`, `REPAIR TABLE`, `RESET`.

- **Replication control statements**. `START SLAVE`, `STOP SLAVE`, `RESET SLAVE`, `CHANGE MASTER TO`.

13.3.4 SAVEPOINT, ROLLBACK TO SAVEPOINT, and RELEASE SAVEPOINT Syntax

```
SAVEPOINT identifier
```

```
ROLLBACK [WORK] TO [SAVEPOINT] identifier
RELEASE SAVEPOINT identifier
```

InnoDB supports the SQL statements SAVEPOINT, ROLLBACK TO SAVEPOINT, RELEASE SAVEPOINT and the optional WORK keyword for ROLLBACK.

The SAVEPOINT statement sets a named transaction savepoint with a name of *identifier*. If the current transaction has a savepoint with the same name, the old savepoint is deleted and a new one is set.

The ROLLBACK TO SAVEPOINT statement rolls back a transaction to the named savepoint without terminating the transaction. Modifications that the current transaction made to rows after the savepoint was set are undone in the rollback, but InnoDB does *not* release the row locks that were stored in memory after the savepoint. (For a new inserted row, the lock information is carried by the transaction ID stored in the row; the lock is not separately stored in memory. In this case, the row lock is released in the undo.) Savepoints that were set at a later time than the named savepoint are deleted.

If the ROLLBACK TO SAVEPOINT statement returns the following error, it means that no savepoint with the specified name exists:

```
ERROR 1305 (42000): SAVEPOINT identifier does not exist
```

The RELEASE SAVEPOINT statement removes the named savepoint from the set of savepoints of the current transaction. No commit or rollback occurs. It is an error if the savepoint does not exist.

All savepoints of the current transaction are deleted if you execute a COMMIT, or a ROLLBACK that does not name a savepoint.

A new savepoint level is created when a stored function is invoked or a trigger is activated. The savepoints on previous levels become unavailable and thus do not conflict with savepoints on the new level. When the function or trigger terminates, any savepoints it created are released and the previous savepoint level is restored.

13.3.5 LOCK TABLES and UNLOCK TABLES Syntax

```
LOCK TABLES
    tbl_name [[AS] alias] lock_type
    [, tbl_name [[AS] alias] lock_type] ...

lock_type:
    READ [LOCAL]
  | [LOW_PRIORITY] WRITE

UNLOCK TABLES
```

MySQL enables client sessions to acquire table locks explicitly for the purpose of cooperating with other sessions for access to tables, or to prevent other sessions from modifying tables during periods when a session requires exclusive access to them. A session can acquire or release locks only for itself. One session cannot acquire locks for another session or release locks held by another session.

Locks may be used to emulate transactions or to get more speed when updating tables. This is explained in more detail later in this section.

LOCK TABLES explicitly acquires table locks for the current client session. Table locks can be acquired for base tables or views. You must have the LOCK TABLES privilege, and the SELECT privilege for each object to be locked.

For view locking, LOCK TABLES adds all base tables used in the view to the set of tables to be locked and locks them automatically. If you lock a table explicitly with LOCK TABLES, any tables used in triggers are also locked implicitly, as described in Section 13.3.5.2, “LOCK TABLES and Triggers”.

`UNLOCK TABLES` explicitly releases any table locks held by the current session. `LOCK TABLES` implicitly releases any table locks held by the current session before acquiring new locks.

Another use for `UNLOCK TABLES` is to release the global read lock acquired with the `FLUSH TABLES WITH READ LOCK` statement, which enables you to lock all tables in all databases. See Section 13.7.6.3, “FLUSH Syntax”. (This is a very convenient way to get backups if you have a file system such as Veritas that can take snapshots in time.)

A table lock protects only against inappropriate reads or writes by other sessions. A session holding a `WRITE` lock can perform table-level operations such as `DROP TABLE` or `TRUNCATE TABLE`. For sessions holding a `READ` lock, `DROP TABLE` and `TRUNCATE TABLE` operations are not permitted.

The following discussion applies only to non-`TEMPORARY` tables. `LOCK TABLES` is permitted (but ignored) for a `TEMPORARY` table. The table can be accessed freely by the session within which it was created, regardless of what other locking may be in effect. No lock is necessary because no other session can see the table.

For information about other conditions on the use of `LOCK TABLES` and statements that cannot be used while `LOCK TABLES` is in effect, see Section 13.3.5.3, “Table-Locking Restrictions and Conditions”

Rules for Lock Acquisition

To acquire table locks within the current session, use the `LOCK TABLES` statement. The following lock types are available:

`READ [LOCAL]` lock:

- The session that holds the lock can read the table (but not write it).
- Multiple sessions can acquire a `READ` lock for the table at the same time.
- Other sessions can read the table without explicitly acquiring a `READ` lock.
- The `LOCAL` modifier enables nonconflicting `INSERT` statements (concurrent inserts) by other sessions to execute while the lock is held. (See Section 8.11.3, “Concurrent Inserts”.) However, `READ LOCAL` cannot be used if you are going to manipulate the database using processes external to the server while you hold the lock. For `InnoDB` tables, `READ LOCAL` is the same as `READ`.

`[LOW_PRIORITY] WRITE` lock:

- The session that holds the lock can read and write the table.
- Only the session that holds the lock can access the table. No other session can access it until the lock is released.
- Lock requests for the table by other sessions block while the `WRITE` lock is held.
- The `LOW_PRIORITY` modifier has no effect. In previous versions of MySQL, it affected locking behavior, but this is no longer true. It is now deprecated and its use produces a warning. Use `WRITE` without `LOW_PRIORITY` instead.

If the `LOCK TABLES` statement must wait due to locks held by other sessions on any of the tables, it blocks until all locks can be acquired.

A session that requires locks must acquire all the locks that it needs in a single `LOCK TABLES` statement. While the locks thus obtained are held, the session can access only the locked tables. For example, in the following sequence of statements, an error occurs for the attempt to access `t2` because it was not locked in the `LOCK TABLES` statement:

```
mysql> LOCK TABLES t1 READ;
mysql> SELECT COUNT(*) FROM t1;
```

```
+----------+
| COUNT(*) |
+----------+
|        3 |
+----------+
mysql> SELECT COUNT(*) FROM t2;
ERROR 1100 (HY000): Table 't2' was not locked with LOCK TABLES
```

Tables in the `INFORMATION_SCHEMA` database are an exception. They can be accessed without being locked explicitly even while a session holds table locks obtained with `LOCK TABLES`.

You cannot refer to a locked table multiple times in a single query using the same name. Use aliases instead, and obtain a separate lock for the table and each alias:

```
mysql> LOCK TABLE t WRITE, t AS t1 READ;
mysql> INSERT INTO t SELECT * FROM t;
ERROR 1100: Table 't' was not locked with LOCK TABLES
mysql> INSERT INTO t SELECT * FROM t AS t1;
```

The error occurs for the first `INSERT` because there are two references to the same name for a locked table. The second `INSERT` succeeds because the references to the table use different names.

If your statements refer to a table by means of an alias, you must lock the table using that same alias. It does not work to lock the table without specifying the alias:

```
mysql> LOCK TABLE t READ;
mysql> SELECT * FROM t AS myalias;
ERROR 1100: Table 'myalias' was not locked with LOCK TABLES
```

Conversely, if you lock a table using an alias, you must refer to it in your statements using that alias:

```
mysql> LOCK TABLE t AS myalias READ;
mysql> SELECT * FROM t;
ERROR 1100: Table 't' was not locked with LOCK TABLES
mysql> SELECT * FROM t AS myalias;
```

`WRITE` locks normally have higher priority than `READ` locks to ensure that updates are processed as soon as possible. This means that if one session obtains a `READ` lock and then another session requests a `WRITE` lock, subsequent `READ` lock requests wait until the session that requested the `WRITE` lock has obtained the lock and released it.

`LOCK TABLES` acquires locks as follows:

1. Sort all tables to be locked in an internally defined order. From the user standpoint, this order is undefined.
2. If a table is to be locked with a read and a write lock, put the write lock request before the read lock request.
3. Lock one table at a time until the session gets all locks.

This policy ensures that table locking is deadlock free.

Note

`LOCK TABLES` or `UNLOCK TABLES`, when applied to a partitioned table, always locks or unlocks the entire table; these statements do not support partition lock pruning. See Partitioning and Locking.

Rules for Lock Release

When the table locks held by a session are released, they are all released at the same time. A session can release its locks explicitly, or locks may be released implicitly under certain conditions.

- A session can release its locks explicitly with `UNLOCK TABLES`.
- If a session issues a `LOCK TABLES` statement to acquire a lock while already holding locks, its existing locks are released implicitly before the new locks are granted.
- If a session begins a transaction (for example, with `START TRANSACTION`), an implicit `UNLOCK TABLES` is performed, which causes existing locks to be released. (For additional information about the interaction between table locking and transactions, see Section 13.3.5.1, “Interaction of Table Locking and Transactions”.)

If the connection for a client session terminates, whether normally or abnormally, the server implicitly releases all table locks held by the session (transactional and nontransactional). If the client reconnects, the locks will no longer be in effect. In addition, if the client had an active transaction, the server rolls back the transaction upon disconnect, and if reconnect occurs, the new session begins with autocommit enabled. For this reason, clients may wish to disable auto-reconnect. With auto-reconnect in effect, the client is not notified if reconnect occurs but any table locks or current transaction will have been lost. With auto-reconnect disabled, if the connection drops, an error occurs for the next statement issued. The client can detect the error and take appropriate action such as reacquiring the locks or redoing the transaction. See Section 27.7.19, “Controlling Automatic Reconnection Behavior”.

Note

If you use `ALTER TABLE` on a locked table, it may become unlocked. For example, if you attempt a second `ALTER TABLE` operation, the result may be an error `Table 'tbl_name' was not locked with LOCK TABLES`. To handle this, lock the table again prior to the second alteration. See also Section B.5.6.1, “Problems with ALTER TABLE”.

13.3.5.1 Interaction of Table Locking and Transactions

`LOCK TABLES` and `UNLOCK TABLES` interact with the use of transactions as follows:

- `LOCK TABLES` is not transaction-safe and implicitly commits any active transaction before attempting to lock the tables.
- `UNLOCK TABLES` implicitly commits any active transaction, but only if `LOCK TABLES` has been used to acquire table locks. For example, in the following set of statements, `UNLOCK TABLES` releases the global read lock but does not commit the transaction because no table locks are in effect:

```
FLUSH TABLES WITH READ LOCK;
START TRANSACTION;
SELECT ... ;
UNLOCK TABLES;
```

- Beginning a transaction (for example, with `START TRANSACTION`) implicitly commits any current transaction and releases existing table locks.
- `FLUSH TABLES WITH READ LOCK` acquires a global read lock and not table locks, so it is not subject to the same behavior as `LOCK TABLES` and `UNLOCK TABLES` with respect to table locking and implicit commits. For example, `START TRANSACTION` does not release the global read lock. See Section 13.7.6.3, “FLUSH Syntax”.
- Other statements that implicitly cause transactions to be committed do not release existing table locks. For a list of such statements, see Section 13.3.3, “Statements That Cause an Implicit Commit”.
- The correct way to use `LOCK TABLES` and `UNLOCK TABLES` with transactional tables, such as `InnoDB` tables, is to begin a transaction with `SET autocommit = 0` (not `START TRANSACTION`) followed by `LOCK TABLES`, and to not call `UNLOCK TABLES` until you commit the transaction explicitly. For example, if you need to write to table `t1` and read from table `t2`, you can do this:

```
SET autocommit=0;
LOCK TABLES t1 WRITE, t2 READ, ...;
... do something with tables t1 and t2 here ...
COMMIT;
UNLOCK TABLES;
```

When you call `LOCK TABLES`, `InnoDB` internally takes its own table lock, and MySQL takes its own table lock. `InnoDB` releases its internal table lock at the next commit, but for MySQL to release its table lock, you have to call `UNLOCK TABLES`. You should not have `autocommit = 1`, because then `InnoDB` releases its internal table lock immediately after the call of `LOCK TABLES`, and deadlocks can very easily happen. `InnoDB` does not acquire the internal table lock at all if `autocommit = 1`, to help old applications avoid unnecessary deadlocks.

- `ROLLBACK` does not release table locks.

13.3.5.2 LOCK TABLES and Triggers

If you lock a table explicitly with `LOCK TABLES`, any tables used in triggers are also locked implicitly:

- The locks are taken as the same time as those acquired explicitly with the `LOCK TABLES` statement.
- The lock on a table used in a trigger depends on whether the table is used only for reading. If so, a read lock suffices. Otherwise, a write lock is used.
- If a table is locked explicitly for reading with `LOCK TABLES`, but needs to be locked for writing because it might be modified within a trigger, a write lock is taken rather than a read lock. (That is, an implicit write lock needed due to the table's appearance within a trigger causes an explicit read lock request for the table to be converted to a write lock request.)

Suppose that you lock two tables, `t1` and `t2`, using this statement:

```
LOCK TABLES t1 WRITE, t2 READ;
```

If `t1` or `t2` have any triggers, tables used within the triggers will also be locked. Suppose that `t1` has a trigger defined like this:

```
CREATE TRIGGER t1_a_ins AFTER INSERT ON t1 FOR EACH ROW
BEGIN
  UPDATE t4 SET count = count+1
      WHERE id = NEW.id AND EXISTS (SELECT a FROM t3);
  INSERT INTO t2 VALUES(1, 2);
END;
```

The result of the `LOCK TABLES` statement is that `t1` and `t2` are locked because they appear in the statement, and `t3` and `t4` are locked because they are used within the trigger:

- `t1` is locked for writing per the `WRITE` lock request.
- `t2` is locked for writing, even though the request is for a `READ` lock. This occurs because `t2` is inserted into within the trigger, so the `READ` request is converted to a `WRITE` request.
- `t3` is locked for reading because it is only read from within the trigger.
- `t4` is locked for writing because it might be updated within the trigger.

13.3.5.3 Table-Locking Restrictions and Conditions

You can safely use `KILL` to terminate a session that is waiting for a table lock. See Section 13.7.6.4, “KILL Syntax”.

`LOCK TABLES` and `UNLOCK TABLES` cannot be used within stored programs.

Tables in the `performance_schema` database cannot be locked with `LOCK TABLES`, except the `setup_xxx` tables.

The following statements are prohibited while a `LOCK TABLES` statement is in effect: `CREATE TABLE`, `CREATE TABLE ... LIKE`, `CREATE VIEW`, `DROP VIEW`, and DDL statements on stored functions and procedures and events.

For some operations, system tables in the `mysql` database must be accessed. For example, the `HELP` statement requires the contents of the server-side help tables, and `CONVERT_TZ()` might need to read the time zone tables. The server implicitly locks the system tables for reading as necessary so that you need not lock them explicitly. These tables are treated as just described:

```
mysql.help_category
mysql.help_keyword
mysql.help_relation
mysql.help_topic
mysql.proc
mysql.time_zone
mysql.time_zone_leap_second
mysql.time_zone_name
mysql.time_zone_transition
mysql.time_zone_transition_type
```

If you want to explicitly place a `WRITE` lock on any of those tables with a `LOCK TABLES` statement, the table must be the only one locked; no other table can be locked with the same statement.

Normally, you do not need to lock tables, because all single `UPDATE` statements are atomic; no other session can interfere with any other currently executing SQL statement. However, there are a few cases when locking tables may provide an advantage:

- If you are going to run many operations on a set of `MyISAM` tables, it is much faster to lock the tables you are going to use. Locking `MyISAM` tables speeds up inserting, updating, or deleting on them because MySQL does not flush the key cache for the locked tables until `UNLOCK TABLES` is called. Normally, the key cache is flushed after each SQL statement.

  The downside to locking the tables is that no session can update a `READ`-locked table (including the one holding the lock) and no session can access a `WRITE`-locked table other than the one holding the lock.

- If you are using tables for a nontransactional storage engine, you must use `LOCK TABLES` if you want to ensure that no other session modifies the tables between a `SELECT` and an `UPDATE`. The example shown here requires `LOCK TABLES` to execute safely:

  ```
  LOCK TABLES trans READ, customer WRITE;
  SELECT SUM(value) FROM trans WHERE customer_id=some_id;
  UPDATE customer
    SET total_value=sum_from_previous_statement
    WHERE customer_id=some_id;
  UNLOCK TABLES;
  ```

  Without `LOCK TABLES`, it is possible that another session might insert a new row in the `trans` table between execution of the `SELECT` and `UPDATE` statements.

You can avoid using `LOCK TABLES` in many cases by using relative updates (`UPDATE customer SET value=value+new_value`) or the `LAST_INSERT_ID()` function.

You can also avoid locking tables in some cases by using the user-level advisory lock functions `GET_LOCK()` and `RELEASE_LOCK()`. These locks are saved in a hash table in the server and implemented with `pthread_mutex_lock()` and `pthread_mutex_unlock()` for high speed. See Section 12.20, “Miscellaneous Functions”.

See Section 8.11.1, “Internal Locking Methods”, for more information on locking policy.

13.3.6 SET TRANSACTION Syntax

```
SET [GLOBAL | SESSION] TRANSACTION
    transaction_characteristic [, transaction_characteristic] ...

transaction_characteristic:
    ISOLATION LEVEL level
  | READ WRITE
  | READ ONLY

level:
     REPEATABLE READ
   | READ COMMITTED
   | READ UNCOMMITTED
   | SERIALIZABLE
```

This statement specifies transaction characteristics. It takes a list of one or more characteristic values separated by commas. These characteristics set the transaction isolation level or access mode. The isolation level is used for operations on `InnoDB` tables. The access mode may be specified as to whether transactions operate in read/write or read-only mode.

In addition, `SET TRANSACTION` can include an optional `GLOBAL` or `SESSION` keyword to indicate the scope of the statement.

Scope of Transaction Characteristics

You can set transaction characteristics globally, for the current session, or for the next transaction:

- With the `GLOBAL` keyword, the statement applies globally for all subsequent sessions. Existing sessions are unaffected.
- With the `SESSION` keyword, the statement applies to all subsequent transactions performed within the current session.
- Without any `SESSION` or `GLOBAL` keyword, the statement applies to the next (not started) transaction performed within the current session. Subsequent transactions revert to using the `SESSION` isolation level.

A global change to transaction characteristics requires the `SUPER` privilege. Any session is free to change its session characteristics (even in the middle of a transaction), or the characteristics for its next transaction.

`SET TRANSACTION` without `GLOBAL` or `SESSION` is not permitted while there is an active transaction:

```
mysql> START TRANSACTION;
Query OK, 0 rows affected (0.02 sec)

mysql> SET TRANSACTION ISOLATION LEVEL SERIALIZABLE;
ERROR 1568 (25001): Transaction characteristics can't be changed
while a transaction is in progress
```

To set the global default isolation level at server startup, use the `--transaction-isolation=level` option to `mysqld` on the command line or in an option file. Values of *`level`* for this option use dashes rather than spaces, so the permissible values are `READ-UNCOMMITTED`, `READ-COMMITTED`, `REPEATABLE-READ`, or `SERIALIZABLE`. For example, to set the default isolation level to `REPEATABLE READ`, use these lines in the `[mysqld]` section of an option file:

```
[mysqld]
transaction-isolation = REPEATABLE-READ
```

It is possible to check or set the global and session transaction isolation levels at runtime by using the `tx_isolation` system variable:

```
SELECT @@GLOBAL.tx_isolation, @@tx_isolation;
SET GLOBAL tx_isolation='REPEATABLE-READ';
SET SESSION tx_isolation='SERIALIZABLE';
```

Similarly, to set the transaction access mode at server startup or at runtime, use the `--transaction-read-only` option or `tx_read_only` system variable. By default, these are `OFF` (the mode is read/write) but can be set to `ON` for a default mode of read only.

Setting the global or session value of `tx_isolation` or `tx_read_only` is equivalent to setting the isolation level or access mode with `SET GLOBAL TRANSACTION` or `SET SESSION TRANSACTION`.

Transaction Isolation Levels

For information about transaction isolation levels, see Section 15.5.2.1, “Transaction Isolation Levels”.

Transaction Access Mode

The transaction access mode may be specified with `SET TRANSACTION`. By default, a transaction takes place in read/write mode, with both reads and writes permitted to tables used in the transaction. This mode may be specified explicitly using an access mode of `READ WRITE`.

If the transaction access mode is set to `READ ONLY`, changes to tables are prohibited. This may enable storage engines to make performance improvements that are possible when writes are not permitted.

It is not permitted to specify both `READ WRITE` and `READ ONLY` in the same statement.

In read-only mode, it remains possible to change tables created with the `TEMPORARY` keyword using DML statements. Changes made with DDL statements are not permitted, just as with permanent tables.

The `READ WRITE` and `READ ONLY` access modes also may be specified for an individual transaction using the `START TRANSACTION` statement.

13.3.7 XA Transactions

Support for XA transactions is available for the `InnoDB` storage engine. The MySQL XA implementation is based on the X/Open CAE document *Distributed Transaction Processing: The XA Specification*. This document is published by The Open Group and available at http://www.opengroup.org/public/pubs/catalog/c193.htm. Limitations of the current XA implementation are described in Section C.6, “Restrictions on XA Transactions”.

On the client side, there are no special requirements. The XA interface to a MySQL server consists of SQL statements that begin with the `XA` keyword. MySQL client programs must be able to send SQL statements and to understand the semantics of the XA statement interface. They do not need be linked against a recent client library. Older client libraries also will work.

Among the MySQL Connectors, MySQL Connector/J 5.0.0 and higher supports XA directly, by means of a class interface that handles the XA SQL statement interface for you.

XA supports distributed transactions, that is, the ability to permit multiple separate transactional resources to participate in a global transaction. Transactional resources often are RDBMSs but may be other kinds of resources.

A global transaction involves several actions that are transactional in themselves, but that all must either complete successfully as a group, or all be rolled back as a group. In essence, this extends ACID properties “up a level” so that multiple ACID transactions can be executed in concert as components of a global operation that also has ACID properties. (As with nondistributed transactions, `SERIALIZABLE` may be preferred if your applications are sensitive to read phenomena. `REPEATABLE READ` may not be sufficient for distributed transactions.)

Some examples of distributed transactions:

- An application may act as an integration tool that combines a messaging service with an RDBMS. The application makes sure that transactions dealing with message sending, retrieval, and processing that also involve a transactional database all happen in a global transaction. You can think of this as “transactional email.”
- An application performs actions that involve different database servers, such as a MySQL server and an Oracle server (or multiple MySQL servers), where actions that involve multiple servers must happen as part of a global transaction, rather than as separate transactions local to each server.
- A bank keeps account information in an RDBMS and distributes and receives money through automated teller machines (ATMs). It is necessary to ensure that ATM actions are correctly reflected in the accounts, but this cannot be done with the RDBMS alone. A global transaction manager integrates the ATM and database resources to ensure overall consistency of financial transactions.

Applications that use global transactions involve one or more Resource Managers and a Transaction Manager:

- A Resource Manager (RM) provides access to transactional resources. A database server is one kind of resource manager. It must be possible to either commit or roll back transactions managed by the RM.
- A Transaction Manager (TM) coordinates the transactions that are part of a global transaction. It communicates with the RMs that handle each of these transactions. The individual transactions within a global transaction are “branches” of the global transaction. Global transactions and their branches are identified by a naming scheme described later.

The MySQL implementation of XA enables a MySQL server to act as a Resource Manager that handles XA transactions within a global transaction. A client program that connects to the MySQL server acts as the Transaction Manager.

To carry out a global transaction, it is necessary to know which components are involved, and bring each component to a point when it can be committed or rolled back. Depending on what each component reports about its ability to succeed, they must all commit or roll back as an atomic group. That is, either all components must commit, or all components must roll back. To manage a global transaction, it is necessary to take into account that any component or the connecting network might fail.

The process for executing a global transaction uses two-phase commit (2PC). This takes place after the actions performed by the branches of the global transaction have been executed.

1. In the first phase, all branches are prepared. That is, they are told by the TM to get ready to commit. Typically, this means each RM that manages a branch records the actions for the branch in stable storage. The branches indicate whether they are able to do this, and these results are used for the second phase.
2. In the second phase, the TM tells the RMs whether to commit or roll back. If all branches indicated when they were prepared that they will be able to commit, all branches are told to commit. If any branch indicated when it was prepared that it will not be able to commit, all branches are told to roll back.

In some cases, a global transaction might use one-phase commit (1PC). For example, when a Transaction Manager finds that a global transaction consists of only one transactional resource (that is, a single branch), that resource can be told to prepare and commit at the same time.

13.3.7.1 XA Transaction SQL Syntax

To perform XA transactions in MySQL, use the following statements:

```
XA {START|BEGIN} xid [JOIN|RESUME]

XA END xid [SUSPEND [FOR MIGRATE]]

XA PREPARE xid

XA COMMIT xid [ONE PHASE]

XA ROLLBACK xid

XA RECOVER [CONVERT XID]
```

For `XA START`, the `JOIN` and `RESUME` clauses are not supported.

For `XA END` the `SUSPEND [FOR MIGRATE]` clause is not supported.

Each XA statement begins with the `XA` keyword, and most of them require an *`xid`* value. An *`xid`* is an XA transaction identifier. It indicates which transaction the statement applies to. *`xid`* values are supplied by the client, or generated by the MySQL server. An *`xid`* value has from one to three parts:

```
xid: gtrid [, bqual [, formatID ]]
```

`gtrid` is a global transaction identifier, *`bqual`* is a branch qualifier, and *`formatID`* is a number that identifies the format used by the *`gtrid`* and *`bqual`* values. As indicated by the syntax, *`bqual`* and *`formatID`* are optional. The default *`bqual`* value is `''` if not given. The default *`formatID`* value is 1 if not given.

`gtrid` and *`bqual`* must be string literals, each up to 64 bytes (not characters) long. *`gtrid`* and *`bqual`* can be specified in several ways. You can use a quoted string (`'ab'`), hex string (`X'6162'`, `0x6162`), or bit value (`b'nnnn'`).

`formatID` is an unsigned integer.

The *`gtrid`* and *`bqual`* values are interpreted in bytes by the MySQL server's underlying XA support routines. However, while an SQL statement containing an XA statement is being parsed, the server works with some specific character set. To be safe, write *`gtrid`* and *`bqual`* as hex strings.

`xid` values typically are generated by the Transaction Manager. Values generated by one TM must be different from values generated by other TMs. A given TM must be able to recognize its own *`xid`* values in a list of values returned by the `XA RECOVER` statement.

For `XA START` *`xid`* starts an XA transaction with the given *`xid`* value. Each XA transaction must have a unique *`xid`* value, so the value must not currently be used by another XA transaction. Uniqueness is assessed using the *`gtrid`* and *`bqual`* values. All following XA statements for the XA transaction must be specified using the same *`xid`* value as that given in the `XA START` statement. If you use any of those statements but specify an *`xid`* value that does not correspond to some existing XA transaction, an error occurs.

One or more XA transactions can be part of the same global transaction. All XA transactions within a given global transaction must use the same *`gtrid`* value in the *`xid`* value. For this reason, *`gtrid`* values must be globally unique so that there is no ambiguity about which global transaction a given XA transaction is part of. The *`bqual`* part of the *`xid`* value must be different for each XA transaction within a global transaction. (The requirement that *`bqual`* values be different is a limitation of the current MySQL XA implementation. It is not part of the XA specification.)

The `XA RECOVER` statement returns information for those XA transactions on the MySQL server that are in the `PREPARED` state. (See Section 13.3.7.2, “XA Transaction States”.) The output includes a row for each such XA transaction on the server, regardless of which client started it.

As of MySQL 8.0.2, `XA RECOVER` requires the `XA_RECOVER_ADMIN` privilege. This privilege requirement prevents users from discovering the XID values for outstanding prepared XA transactions

other than their own. It does not affect normal commit or rollback of an XA transaction because the user who started it knows its XID.

`XA RECOVER` output rows look like this (for an example *`xid`* value consisting of the parts `'abc'`, `'def'`, and `7`):

```
mysql> XA RECOVER;
+----------+--------------+--------------+--------+
| formatID | gtrid_length | bqual_length | data   |
+----------+--------------+--------------+--------+
|        7 |            3 |            3 | abcdef |
+----------+--------------+--------------+--------+
```

The output columns have the following meanings:

- `formatID` is the *`formatID`* part of the transaction *`xid`*
- `gtrid_length` is the length in bytes of the *`gtrid`* part of the *`xid`*
- `bqual_length` is the length in bytes of the *`bqual`* part of the *`xid`*
- `data` is the concatenation of the *`gtrid`* and *`bqual`* parts of the *`xid`*

XID values may contain nonprintable characters. `XA RECOVER` permits an optional `CONVERT XID` clause so that clients can request XID values in hexadecimal.

13.3.7.2 XA Transaction States

An XA transaction progresses through the following states:

1. Use `XA START` to start an XA transaction and put it in the `ACTIVE` state.
2. For an `ACTIVE` XA transaction, issue the SQL statements that make up the transaction, and then issue an `XA END` statement. `XA END` puts the transaction in the `IDLE` state.
3. For an `IDLE` XA transaction, you can issue either an `XA PREPARE` statement or an `XA COMMIT ... ONE PHASE` statement:
   - `XA PREPARE` puts the transaction in the `PREPARED` state. An `XA RECOVER` statement at this point will include the transaction's *`xid`* value in its output, because `XA RECOVER` lists all XA transactions that are in the `PREPARED` state.
   - `XA COMMIT ... ONE PHASE` prepares and commits the transaction. The *`xid`* value will not be listed by `XA RECOVER` because the transaction terminates.
4. For a `PREPARED` XA transaction, you can issue an `XA COMMIT` statement to commit and terminate the transaction, or `XA ROLLBACK` to roll back and terminate the transaction.

Here is a simple XA transaction that inserts a row into a table as part of a global transaction:

```
mysql> XA START 'xatest';
Query OK, 0 rows affected (0.00 sec)

mysql> INSERT INTO mytable (i) VALUES(10);
Query OK, 1 row affected (0.04 sec)

mysql> XA END 'xatest';
Query OK, 0 rows affected (0.00 sec)

mysql> XA PREPARE 'xatest';
Query OK, 0 rows affected (0.00 sec)

mysql> XA COMMIT 'xatest';
```

```
Query OK, 0 rows affected (0.00 sec)
```

Within the context of a given client connection, XA transactions and local (non-XA) transactions are mutually exclusive. For example, if `XA START` has been issued to begin an XA transaction, a local transaction cannot be started until the XA transaction has been committed or rolled back. Conversely, if a local transaction has been started with `START TRANSACTION`, no XA statements can be used until the transaction has been committed or rolled back.

If an XA transaction is in the `ACTIVE` state, you cannot issue any statements that cause an implicit commit. That would violate the XA contract because you could not roll back the XA transaction. You will receive the following error if you try to execute such a statement:

```
ERROR 1399 (XAE07): XAER_RMFAIL: The command cannot be executed
when global transaction is in the ACTIVE state
```

Statements to which the preceding remark applies are listed at Section 13.3.3, “Statements That Cause an Implicit Commit”.

13.4 Replication Statements

Replication can be controlled through the SQL interface using the statements described in this section. Statements are split into a group which controls master servers, a group which controls slave servers, and a group which can be applied to any replication servers.

13.4.1 SQL Statements for Controlling Master Servers

This section discusses statements for managing master replication servers. Section 13.4.2, “SQL Statements for Controlling Slave Servers”, discusses statements for managing slave servers.

In addition to the statements described here, the following `SHOW` statements are used with master servers in replication. For information about these statements, see Section 13.7.5, “SHOW Syntax”.

- `SHOW BINARY LOGS`
- `SHOW BINLOG EVENTS`
- `SHOW MASTER STATUS`
- `SHOW SLAVE HOSTS`

13.4.1.1 PURGE BINARY LOGS Syntax

```
PURGE { BINARY | MASTER } LOGS
    { TO 'log_name' | BEFORE datetime_expr }
```

The binary log is a set of files that contain information about data modifications made by the MySQL server. The log consists of a set of binary log files, plus an index file (see Section 5.4.4, “The Binary Log”).

The `PURGE BINARY LOGS` statement deletes all the binary log files listed in the log index file prior to the specified log file name or date. `BINARY` and `MASTER` are synonyms. Deleted log files also are removed from the list recorded in the index file, so that the given log file becomes the first in the list.

This statement has no effect if the server was not started with the `--log-bin` option to enable binary logging.

Examples:

```
PURGE BINARY LOGS TO 'mysql-bin.010';
PURGE BINARY LOGS BEFORE '2008-04-02 22:46:26';
```

The BEFORE variant's *datetime_expr* argument should evaluate to a DATETIME value (a value in 'YYYY-MM-DD hh:mm:ss' format).

This statement is safe to run while slaves are replicating. You need not stop them. If you have an active slave that currently is reading one of the log files you are trying to delete, this statement does nothing and it fails with an error in such cases. However, if a slave is not connected and you happen to purge one of the log files it has yet to read, the slave will be unable to replicate after it reconnects.

To safely purge binary log files, follow this procedure:

1. On each slave server, use SHOW SLAVE STATUS to check which log file it is reading.
2. Obtain a listing of the binary log files on the master server with SHOW BINARY LOGS.
3. Determine the earliest log file among all the slaves. This is the target file. If all the slaves are up to date, this is the last log file on the list.
4. Make a backup of all the log files you are about to delete. (This step is optional, but always advisable.)
5. Purge all log files up to but not including the target file.

You can also set the expire_logs_days system variable to expire binary log files automatically after a given number of days (see Section 5.1.5, “Server System Variables”). If you are using replication, you should set the variable no lower than the maximum number of days your slaves might lag behind the master.

PURGE BINARY LOGS TO and PURGE BINARY LOGS BEFORE both fail with an error when binary log files listed in the .index file had been removed from the system by some other means (such as using rm on Linux). (Bug #18199, Bug #18453) To handle such errors, edit the .index file (which is a simple text file) manually to ensure that it lists only the binary log files that are actually present, then run again the PURGE BINARY LOGS statement that failed.

13.4.1.2 RESET MASTER Syntax

```
RESET MASTER [TO binary_log_file_index_number]
```

RESET MASTER enables you to delete any binary log files and their related binary log index file, returning the master to its state before binary logging was started.

Warning

Use this statement with caution to ensure you do not lose binary log file data.

Issuing RESET MASTER without the optional TO clause deletes all binary log files listed in the index file, resets the binary log index file to be empty, and creates a new binary log file starting at 1. Use the optional TO clause to start the binary log file index from a number other than 1 after the reset. Issuing RESET MASTER also clears the values of the gtid_purged system variable and the gtid_executed system variable; that is, issuing this statement sets each of these values to an empty string (''). This statement also clears the mysql.gtid_executed table (see mysql.gtid_executed Table).

Using RESET MASTER with the TO clause to specify a binary log file index number to start from simplifies failover by providing a single statement alternative to the FLUSH BINARY LOGS and PURGE BINARY LOGS TO statements.

The following example demonstrates TO clause usage:

```
RESET MASTER TO 1234;

SHOW BINARY LOGS;
+------------------+-----------+
```

```
| Log_name          | File_size |
+-------------------+-----------+
| master-bin.001234 |       154 |
+-------------------+-----------+
```

Important

The effects of `RESET MASTER` without the `TO` clause differ from those of `PURGE BINARY LOGS` in 2 key ways:

1. `RESET MASTER` removes *all* binary log files that are listed in the index file, leaving only a single, empty binary log file with a numeric suffix of `.000001`, whereas the numbering is not reset by `PURGE BINARY LOGS`.
2. `RESET MASTER` is *not* intended to be used while any replication slaves are running. The behavior of `RESET MASTER` when used while slaves are running is undefined (and thus unsupported), whereas `PURGE BINARY LOGS` may be safely used while replication slaves are running.

See also Section 13.4.1.1, “PURGE BINARY LOGS Syntax”.

`RESET MASTER` without the `TO` clause can prove useful when you first set up the master and the slave, so that you can verify the setup as follows:

1. Start the master and slave, and start replication (see Section 18.1.2, “Setting Up Binary Log File Position Based Replication”).
2. Execute a few test queries on the master.
3. Check that the queries were replicated to the slave.
4. When replication is running correctly, issue `STOP SLAVE` followed by `RESET SLAVE` on the slave, then verify that no unwanted data from the test queries exists on the slave.
5. Issue `RESET MASTER` on the master to clean up the test queries.

After verifying the setup, resetting the master and slave and ensuring that no unwanted data or binary log files generated by testing remain on the master or slave, you can start the slave and begin replicating.

13.4.1.3 SET sql_log_bin Syntax

```
SET sql_log_bin = {0|1}
```

The `sql_log_bin` variable controls whether logging to the binary log is done. The default value is 1 (do logging). To change logging for the current session, change the session value of this variable. The session user must have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege to set this variable. Set this variable to 0 for a session to temporarily disable binary logging while making changes to the master which you do not want to replicate to the slave.

As of MySQL 5.5, `sql_log_bin` can be set as a global or session variable. Setting `sql_log_bin` globally is only detected when a new session is started. Any sessions previously running are not impacted when setting `sql_log_bin` globally.

Warning

Incorrect use of `sql_log_bin` with a global scope means any changes made in an already running session are *still* being recorded to the binary log and therefore replicated. Exercise extreme caution using `sql_log_bin` with a global scope as the above situation could cause unexpected results including replication failure.

In MySQL 5.7, it is not possible to set `@@session.sql_log_bin` within a transaction or subquery. (Bug #53437)

13.4.2 SQL Statements for Controlling Slave Servers

This section discusses statements for managing slave replication servers. Section 13.4.1, “SQL Statements for Controlling Master Servers”, discusses statements for managing master servers.

In addition to the statements described here, `SHOW SLAVE STATUS` and `SHOW RELAYLOG EVENTS` are also used with replication slaves. For information about these statements, see Section 13.7.5.34, “SHOW SLAVE STATUS Syntax”, and Section 13.7.5.32, “SHOW RELAYLOG EVENTS Syntax”.

13.4.2.1 CHANGE MASTER TO Syntax

```
CHANGE MASTER TO option [, option] ... [ channel_option ]

option:
    MASTER_BIND = 'interface_name'
  | MASTER_HOST = 'host_name'
  | MASTER_USER = 'user_name'
  | MASTER_PASSWORD = 'password'
  | MASTER_PORT = port_num
  | MASTER_CONNECT_RETRY = interval
  | MASTER_RETRY_COUNT = count
  | MASTER_DELAY = interval
  | MASTER_HEARTBEAT_PERIOD = interval
  | MASTER_LOG_FILE = 'master_log_name'
  | MASTER_LOG_POS = master_log_pos
  | MASTER_AUTO_POSITION = {0|1}
  | RELAY_LOG_FILE = 'relay_log_name'
  | RELAY_LOG_POS = relay_log_pos
  | MASTER_SSL = {0|1}
  | MASTER_SSL_CA = 'ca_file_name'
  | MASTER_SSL_CAPATH = 'ca_directory_name'
  | MASTER_SSL_CERT = 'cert_file_name'
  | MASTER_SSL_CRL = 'crl_file_name'
  | MASTER_SSL_CRLPATH = 'crl_directory_name'
  | MASTER_SSL_KEY = 'key_file_name'
  | MASTER_SSL_CIPHER = 'cipher_list'
  | MASTER_SSL_VERIFY_SERVER_CERT = {0|1}
  | MASTER_TLS_VERSION = 'protocol_list'
  | IGNORE_SERVER_IDS = (server_id_list)

channel_option:
    FOR CHANNEL channel

server_id_list:
    [server_id [, server_id] ... ]
```

`CHANGE MASTER TO` changes the parameters that the slave server uses for connecting to the master server, for reading the master binary log, and reading the slave relay log. It also updates the contents of the master info and relay log info repositories (see Section 18.2.4, “Replication Relay and Status Logs”). `CHANGE MASTER TO` requires the `REPLICATION_SLAVE_ADMIN` or `SUPER` privilege.

You can issue `CHANGE MASTER TO` statements on a running slave without first stopping it, depending on the states of the slave SQL thread and slave I/O thread. The rules governing such use are provided later in this section.

When using a multi-threaded slave (in other words `slave_parallel_workers` is greater than 0), stopping the slave can cause “gaps” in the sequence of transactions that have been executed from the relay log, regardless of whether the slave was stopped intentionally or otherwise. When such gaps exist, issuing `CHANGE MASTER TO` fails. The solution in this situation is to issue `START SLAVE UNTIL SQL_AFTER_MTS_GAPS` which ensures that the gaps are closed.

The optional `FOR CHANNEL channel` clause enables you to choose which replication channel the statement applies to. If no clause is set and no extra channels exist, the statement applies to the

default channel and behaves the same as versions of MySQL prior to 5.7.6. Providing a `FOR CHANNEL channel` clause applies the `CHANGE MASTER TO` statement to a specific replication channel, and is used to add a new channel or modify an existing channel. For example, to add a new channel called channel2:

```
CHANGE MASTER TO MASTER_NAME=host1, MASTER_PORT=3002 FOR CHANNEL channel2
```

When using multiple replication channels, if a `CHANGE MASTER TO` statement does not have a channel defined using a `FOR CHANNEL channel` clause an error is generated. See Section 18.2.3, "Replication Channels" for more information.

Options not specified retain their value, except as indicated in the following discussion. Thus, in most cases, there is no need to specify options that do not change.

`MASTER_HOST`, `MASTER_USER`, `MASTER_PASSWORD`, and `MASTER_PORT` provide information to the slave about how to connect to its master:

- `MASTER_HOST` and `MASTER_PORT` are the host name (or IP address) of the master host and its TCP/IP port.

  Note

  Replication cannot use Unix socket files. You must be able to connect to the master MySQL server using TCP/IP.

  If you specify the `MASTER_HOST` or `MASTER_PORT` option, the slave assumes that the master server is different from before (even if the option value is the same as its current value.) In this case, the old values for the master binary log file name and position are considered no longer applicable, so if you do not specify `MASTER_LOG_FILE` and `MASTER_LOG_POS` in the statement, `MASTER_LOG_FILE=''` and `MASTER_LOG_POS=4` are silently appended to it.

  Setting `MASTER_HOST=''` (that is, setting its value explicitly to an empty string) is *not* the same as not setting `MASTER_HOST` at all. Beginning with MySQL 5.5, trying to set `MASTER_HOST` to an empty string fails with an error. Previously, setting `MASTER_HOST` to an empty string caused `START SLAVE` subsequently to fail. (Bug #28796)

  Values used for `MASTER_HOST` and other `CHANGE MASTER TO` options are checked for linefeed (`\n` or `0x0A`) characters; the presence of such characters in these values causes the statement to fail with **`ER_MASTER_INFO`**. (Bug #11758581, Bug #50801)

- `MASTER_USER` and `MASTER_PASSWORD` are the user name and password of the account to use for connecting to the master.

  `MASTER_USER` cannot be made empty; setting `MASTER_USER = ''` or leaving it unset when setting a value for `MASTER_PASSWORD` causes an error (Bug #13427949).

  The password used for a MySQL Replication slave account in a `CHANGE MASTER TO` statement is limited to 32 characters in length; trying to use a password of more than 32 characters causes `CHANGE MASTER TO` to fail.

  Note

  Prior to MySQL 5.7.5, if the password was longer, the statement succeeded but any excess characters were silently truncated.

  The text of a running `CHANGE MASTER TO` statement, including values for `MASTER_USER` and `MASTER_PASSWORD`, can be seen in the output of a concurrent `SHOW PROCESSLIST` statement. (The complete text of a `START SLAVE` statement is also visible to `SHOW PROCESSLIST`.)

The `MASTER_SSL_xxx` options provide information about using SSL for the connection. They correspond to the `--ssl-xxx` options described in Section 6.4.5, "Command Options for Secure

Connections", and Section 18.3.9, "Setting Up Replication to Use Secure Connections". These options can be changed even on slaves that are compiled without SSL support. They are saved to the master info repository, but are ignored if the slave does not have SSL support enabled.

The `MASTER_SSL=1` is prescriptive, not advisory. When given, the slave connection to the master must use SSL or the connection attempt fails.

The `MASTER_TLS_VERSION` option specifies the encryption protocols permitted by the master for slave connections. The value is like that for the `tls_version` system variable: A comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, "Secure Connection Protocols and Ciphers".

`MASTER_CONNECT_RETRY` specifies how many seconds to wait between connect retries. The default is 60.

`MASTER_RETRY_COUNT` limits the *number* of reconnection attempts and updates the value of the `Master_Retry_Count` column in the output of `SHOW SLAVE STATUS`. The default value is 24 * 3600 = 86400. `MASTER_RETRY_COUNT` is intended to replace the older `--master-retry-count` server option, and is now the preferred method for setting this limit. You are encouraged not to rely on `--master-retry-count` in new applications and, when upgrading to MySQL 8.0, to update any existing applications that rely on it, so that they use `CHANGE MASTER TO ... MASTER_RETRY_COUNT` instead.

`MASTER_DELAY` specifies how many seconds behind the master the slave must lag. An event received from the master is not executed until at least *`interval`* seconds later than its execution on the master. The default is 0. An error occurs if *`interval`* is not a nonnegative integer in the range from 0 to $2^{31}-1$. For more information, see Section 18.3.11, "Delayed Replication".

A `CHANGE MASTER TO` statement employing the `MASTER_DELAY` option can be executed on a running slave when the slave SQL thread is stopped.

`MASTER_BIND` is for use on replication slaves having multiple network interfaces, and determines which of the slave's network interfaces is chosen for connecting to the master.

The address configured with this option, if any, can be seen in the `Master_Bind` column of the output from `SHOW SLAVE STATUS`. If you are using slave status log tables (server started with `--master-info-repository=TABLE`), the value can also be seen as the `Master_bind` column of the `mysql.slave_master_info` table.

`MASTER_HEARTBEAT_PERIOD` sets the interval in seconds between replication heartbeats. Whenever the master's binary log is updated with an event, the waiting period for the next heartbeat is reset. *`interval`* is a decimal value having the range 0 to 4294967 seconds and a resolution in milliseconds; the smallest nonzero value is 0.001. Heartbeats are sent by the master only if there are no unsent events in the binary log file for a period longer than *`interval`*.

If you are logging master connection information to tables, `MASTER_HEARTBEAT_PERIOD` can be seen as the value of the `Heartbeat` column of the `mysql.slave_master_info` table.

Setting *`interval`* to 0 disables heartbeats altogether. The default value for *`interval`* is equal to the value of `slave_net_timeout` divided by 2.

Setting `@@global.slave_net_timeout` to a value less than that of the current heartbeat interval results in a warning being issued. The effect of issuing `RESET SLAVE` on the heartbeat interval is to reset it to the default value.

`MASTER_LOG_FILE` and `MASTER_LOG_POS` are the coordinates at which the slave I/O thread should begin reading from the master the next time the thread starts. `RELAY_LOG_FILE` and `RELAY_LOG_POS` are the coordinates at which the slave SQL thread should begin reading from the relay log the next time the thread starts. If you specify either of `MASTER_LOG_FILE` or

`MASTER_LOG_POS`, you cannot specify `RELAY_LOG_FILE` or `RELAY_LOG_POS`. If you specify either of `MASTER_LOG_FILE` or `MASTER_LOG_POS`, you also cannot specify `MASTER_AUTO_POSITION = 1` (described later in this section). If neither of `MASTER_LOG_FILE` or `MASTER_LOG_POS` is specified, the slave uses the last coordinates of the *slave SQL thread* before `CHANGE MASTER TO` was issued. This ensures that there is no discontinuity in replication, even if the slave SQL thread was late compared to the slave I/O thread, when you merely want to change, say, the password to use.

A `CHANGE MASTER TO` statement employing `RELAY_LOG_FILE`, `RELAY_LOG_POS`, or both options can be executed on a running slave when the slave SQL thread is stopped.

If `MASTER_AUTO_POSITION = 1` is used with `CHANGE MASTER TO`, the slave attempts to connect to the master using the GTID-based replication protocol. This option can be employed by `CHANGE MASTER TO` only if both the slave SQL and slave I/O threads are stopped.

When using GTIDs, the slave tells the master which transactions it has already received, executed, or both. To compute this set, it reads the global value of `gtid_executed` and the value of the `Retrieved_gtid_set` column from `SHOW SLAVE STATUS`. Since the GTID of the last transmitted transaction is included in `Retrieved_gtid_set` even if the transaction was only partially transmitted, the last received GTID is subtracted from this set. Thus, the slave computes the following set:

```
UNION(@@global.gtid_executed, Retrieved_gtid_set - last_received_GTID)
```

This set is sent to the master as part of the initial handshake, and the master sends back all transactions that it has executed which are not part of the set. If any of these transactions have been already purged from the master's binary log, the master sends the error **`ER_MASTER_HAS_PURGED_REQUIRED_GTIDS`** to the slave, and replication does not start.

When GTID-based replication is employed, the coordinates represented by `MASTER_LOG_FILE` and `MASTER_LOG_POS` are not used, and global transaction identifiers are used instead. Thus the use of either or both of these options together with `MASTER_AUTO_POSITION` causes an error.

You can see whether replication is running with autopositioning enabled by checking the output of `SHOW SLAVE STATUS`.

`gtid_mode` must also be enabled before issuing `CHANGE MASTER TO ... MASTER_AUTO_POSITION = 1`. Otherwise, the statement fails with an error.

To revert to the older file-based replication protocol after using GTIDs, you can issue a new `CHANGE MASTER TO` statement that specifies `MASTER_AUTO_POSITION = 0`, as well as at least one of `MASTER_LOG_FILE` or `MASTER_LOG_POSITION`.

Relay logs are preserved when neither the slave SQL thread nor the slave I/O thread is stopped; if both threads are stopped, all relay log files are deleted unless you at least one of `RELAY_LOG_FILE` or `RELAY_LOG_POS` is specified.

`RELAY_LOG_FILE` can use either an absolute or relative path, and uses the same base name as `MASTER_LOG_FILE`. (Bug #12190)

`IGNORE_SERVER_IDS` takes a comma-separated list of 0 or more server IDs. Events originating from the corresponding servers are ignored, with the exception of log rotation and deletion events, which are still recorded in the relay log.

In circular replication, the originating server normally acts as the terminator of its own events, so that they are not applied more than once. Thus, this option is useful in circular replication when one of the servers in the circle is removed. Suppose that you have a circular replication setup with 4 servers, having server IDs 1, 2, 3, and 4, and server 3 fails. When bridging the gap by starting replication from server 2 to server 4, you can include `IGNORE_SERVER_IDS = (3)` in the `CHANGE MASTER TO` statement that you issue on server 4 to tell it to use server 2 as its master instead of server 3. Doing so causes it to ignore and not to propagate any statements that originated with the server that is no longer in use.

If a `CHANGE MASTER TO` statement is issued without any `IGNORE_SERVER_IDS` option, any existing list is preserved. To clear the list of ignored servers, it is necessary to use the option with an empty list:

```
CHANGE MASTER TO IGNORE_SERVER_IDS = ();
```

`RESET SLAVE ALL` clears `IGNORE_SERVER_IDS`. Before MySQL 5.7, it did not.

If `IGNORE_SERVER_IDS` contains the server's own ID and the server was started with the `--replicate-same-server-id` option enabled, an error results.

In MySQL 8.0, the master info repository and the output of `SHOW SLAVE STATUS` provide the list of servers that are currently ignored. For more information, see Section 18.2.4.2, “Slave Status Logs”, and Section 13.7.5.34, “SHOW SLAVE STATUS Syntax”.

In MySQL 8.0, invoking `CHANGE MASTER TO` causes the previous values for `MASTER_HOST`, `MASTER_PORT`, `MASTER_LOG_FILE`, and `MASTER_LOG_POS` to be written to the error log, along with other information about the slave's state prior to execution.

In MySQL 8.0, `CHANGE MASTER TO` causes an implicit commit of an ongoing transaction. See Section 13.3.3, “Statements That Cause an Implicit Commit”.

The strict requirement to execute `STOP SLAVE` prior to issuing any `CHANGE MASTER TO` statement (and `START SLAVE` afterward) is removed. Instead of depending on whether the slave is stopped, the behavior of `CHANGE MASTER TO` depends on the states of the slave SQL thread and slave I/O threads; which of these threads is stopped or running now determines the options that can or cannot be used with a `CHANGE MASTER TO` statement at a given point in time. The rules for making this determination are listed here:

- If the SQL thread is stopped, you can execute `CHANGE MASTER TO` using any combination that is otherwise allowed of `RELAY_LOG_FILE`, `RELAY_LOG_POS`, and `MASTER_DELAY` options, even if the slave I/O thread is running. No other options may be used with this statement when the I/O thread is running.
- If the I/O thread is stopped, you can execute `CHANGE MASTER TO` using any of the options for this statement (in any allowed combination) *except* `RELAY_LOG_FILE`, `RELAY_LOG_POS`, or `MASTER_DELAY`, even when the SQL thread is running. These three options may not be used when the I/O thread is running.
- Both the SQL thread and the I/O thread must be stopped before issuing a `CHANGE MASTER TO` statement that employs `MASTER_AUTO_POSITION = 1`.

You can check the current state of the slave SQL and I/O threads using `SHOW SLAVE STATUS`.

For more information, see Section 18.3.8, “Switching Masters During Failover”.

If you are using statement-based replication and temporary tables, it is possible for a `CHANGE MASTER TO` statement following a `STOP SLAVE` statement to leave behind temporary tables on the slave. A warning (**`ER_WARN_OPEN_TEMP_TABLES_MUST_BE_ZERO`**) is now issued whenever this occurs. You can avoid this in such cases by making sure that the value of the `Slave_open_temp_tables` system status variable is equal to 0 prior to executing such a `CHANGE MASTER TO` statement.

`CHANGE MASTER TO` is useful for setting up a slave when you have the snapshot of the master and have recorded the master binary log coordinates corresponding to the time of the snapshot. After loading the snapshot into the slave to synchronize it with the master, you can run `CHANGE MASTER TO MASTER_LOG_FILE='`*`log_name`*`', MASTER_LOG_POS=`*`log_pos`* on the slave to specify the coordinates at which the slave should begin reading the master binary log.

The following example changes the master server the slave uses and establishes the master binary log coordinates from which the slave begins reading. This is used when you want to set up the slave to replicate the master:

```
CHANGE MASTER TO
  MASTER_HOST='master2.mycompany.com',
  MASTER_USER='replication',
  MASTER_PASSWORD='bigs3cret',
  MASTER_PORT=3306,
  MASTER_LOG_FILE='master2-bin.001',
  MASTER_LOG_POS=4,
  MASTER_CONNECT_RETRY=10;
```

The next example shows an operation that is less frequently employed. It is used when the slave has relay log files that you want it to execute again for some reason. To do this, the master need not be reachable. You need only use `CHANGE MASTER TO` and start the SQL thread (`START SLAVE SQL_THREAD`):

```
CHANGE MASTER TO
  RELAY_LOG_FILE='slave-relay-bin.006',
  RELAY_LOG_POS=4025;
```

You can even use the second operation in a nonreplication setup with a standalone, nonslave server for recovery following a crash. Suppose that your server has crashed and you have restored it from a backup. You want to replay the server's own binary log files (not relay log files, but regular binary log files), named (for example) `myhost-bin.*`. First, make a backup copy of these binary log files in some safe place, in case you don't exactly follow the procedure below and accidentally have the server purge the binary log. Use `SET GLOBAL relay_log_purge=0` for additional safety. Then start the server without the `--log-bin` option, Instead, use the `--replicate-same-server-id`, `--relay-log=myhost-bin` (to make the server believe that these regular binary log files are relay log files) and `--skip-slave-start` options. After the server starts, issue these statements:

```
CHANGE MASTER TO
  RELAY_LOG_FILE='myhost-bin.153',
  RELAY_LOG_POS=410,
  MASTER_HOST='some_dummy_string';
START SLAVE SQL_THREAD;
```

The server reads and executes its own binary log files, thus achieving crash recovery. Once the recovery is finished, run `STOP SLAVE`, shut down the server, clear the master info and relay log info repositories, and restart the server with its original options.

Specifying the `MASTER_HOST` option (even with a dummy value) is required to make the server think it is a slave.

The following table shows the maximum permissible length for the string-valued options.

Option	Maximum Length
`MASTER_HOST`	60
`MASTER_USER`	16
`MASTER_PASSWORD`	32
`MASTER_LOG_FILE`	255
`RELAY_LOG_FILE`	255
`MASTER_SSL_CA`	255
`MASTER_SSL_CAPATH`	255
`MASTER_SSL_CERT`	255
`MASTER_SSL_CRL`	255
`MASTER_SSL_CRLPATH`	255
`MASTER_SSL_KEY`	255

Option	Maximum Length
MASTER_SSL_CIPHER	511

13.4.2.2 CHANGE REPLICATION FILTER Syntax

```
CHANGE REPLICATION FILTER filter[, filter]
 [, ...] [FOR CHANNEL channel_name]

filter:
    REPLICATE_DO_DB = (db_list)
  | REPLICATE_IGNORE_DB = (db_list)
  | REPLICATE_DO_TABLE = (tbl_list)
  | REPLICATE_IGNORE_TABLE = (tbl_list)
  | REPLICATE_WILD_DO_TABLE = (wild_tbl_list)
  | REPLICATE_WILD_IGNORE_TABLE = (wild_tbl_list)
  | REPLICATE_REWRITE_DB = (db_pair_list)

db_list:
    db_name[, db_name][, ...]

tbl_list:
    db_name.table_name[, db_table_name][, ...]
wild_tbl_list:
    'db_pattern.table_pattern'[, 'db_pattern.table_pattern'][, ...]

db_pair_list:
    (db_pair)[, (db_pair)][, ...]

db_pair:
    from_db, to_db
```

`CHANGE REPLICATION FILTER` sets one or more replication filtering rules on the slave in the same way as starting the slave `mysqld` with replication filtering options such as `--replicate-do-db` or `--replicate-wild-ignore-table`. Unlike the case with the server options, this statement does not require restarting the server to take effect, only that the slave SQL thread be stopped using `STOP SLAVE SQL_THREAD` first (and restarted with `START SLAVE SQL_THREAD` afterwards). `CHANGE REPLICATION FILTER` requires the `REPLICATION_SLAVE_ADMIN` or `SUPER` privilege. Use the `FOR CHANNEL channel_name` clause to make a replication filter specific to a replication channel, for example on a multi-source replication slave. Filters applied without a specific `FOR CHANNEL` clause are considered global filters, meaning that they are applied to all replication channels.

The following list shows the `CHANGE REPLICATION FILTER` options and how they relate to `--replicate-*` server options:

- `REPLICATE_DO_DB`: Include updates based on database name. Equivalent to `--replicate-do-db`.
- `REPLICATE_IGNORE_DB`: Exclude updates based on database name. Equivalent to `--replicate-ignore-db`.
- `REPLICATE_DO_TABLE`: Include updates based on table name. Equivalent to `--replicate-do-table`.
- `REPLICATE_IGNORE_TABLE`: Exclude updates based on table name. Equivalent to `--replicate-ignore-table`.
- `REPLICATE_WILD_DO_TABLE`: Include updates based on wildcard pattern matching table name. Equivalent to `--replicate-wild-do-table`.
- `REPLICATE_WILD_IGNORE_TABLE`: Exclude updates based on wildcard pattern matching table name. Equivalent to `--replicate-wild-ignore-table`.
- `REPLICATE_REWRITE_DB`: Perform updates on slave after substituting new name on slave for specified database on master. Equivalent to `--replicate-rewrite-db`.

The precise effects of REPLICATE_DO_DB and REPLICATE_IGNORE_DB filters are dependent on whether statement-based or row-based replication is in effect. See Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”, for more information.

Multiple replication filtering rules can be created in a single CHANGE REPLICATION FILTER statement by separating the rules with commas, as shown here:

```
CHANGE REPLICATION FILTER
    REPLICATE_DO_DB = (d1), REPLICATE_IGNORE_DB = (d2);
```

Issuing the statement just shown is equivalent to starting the slave mysqld with the options --replicate-do-db=d1 --replicate-ignore-db=d2.

On a multi-source replication slave, which uses multiple replication channels to process transaction from different sources, use the FOR CHANNEL *channel_name* clause to set a replication filter on a replication channel:

```
CHANGE REPLICATION FILTER REPLICATE_DO_DB = (d1) FOR CHANNEL channel_1;
```

This enables you to create a channel specific replication filter to filter out selected data from a source. When a FOR CHANNEL clause is provided, the replication filter statement acts on that slave replication channel removing any existing replication filter which has the same filter type as the specified replication filters, and replacing them with the specified filter. Filter types not explicitly listed in the statement are not modified. If issued against a slave replication channel which is not configured, the statement fails with an ER_SLAVE_CONFIGURATION error. If issued against Group Replication channels, the statement fails with an ER_SLAVE_CHANNEL_OPERATION_NOT_ALLOWED error.

On a replication slave with multiple replication channels configured, issuing CHANGE REPLICATION FILTER with noFOR CHANNEL clause configures the replication filter for every configured slave replication channel, and for the global replication filters. For every filter type, if the filter type is listed in the statement, then any existing filter rules of that type are replaced by the filter rules specified in the most recently issued statement, otherwise the old value of the filter type is retained. For more information see Section 18.2.5.4, “Replication Channel Based Filters”.

If the same filtering rule is specified multiple times, only the *last* such rule is actually used. For example, the two statements shown here have exactly the same effect, because the first REPLICATE_DO_DB rule in the first statement is ignored:

```
CHANGE REPLICATION FILTER
    REPLICATE_DO_DB = (db1, db2), REPLICATE_DO_DB = (db3, db4);

CHANGE REPLICATION FILTER
    REPLICATE_DO_DB = (db3,db4);
```

Caution

This behavior differs from that of the --replicate-* filter options where specifying the same option multiple times causes the creation of multiple filter rules.

Names of tables and database not containing any special characters need not be quoted. Values used with REPLICATION_WILD_TABLE and REPLICATE_WILD_IGNORE_TABLE are string expressions, possibly containing (special) wildcard characters, and so must be quoted. This is shown in the following example statements:

```
CHANGE REPLICATION FILTER
    REPLICATE_WILD_DO_TABLE = ('db1.old%');

CHANGE REPLICATION FILTER
    REPLICATE_WILD_IGNORE_TABLE = ('db1.new%', 'db2.new%');
```

Values used with REPLICATE_REWRITE_DB represent *pairs* of database names; each such value must be enclosed in parentheses. The following statement rewrites statements occurring on database db1 on the master to database db2 on the slave:

```
CHANGE REPLICATION FILTER REPLICATE_REWRITE_DB = ((db1, db2));
```

The statement just shown contains two sets of parentheses, one enclosing the pair of database names, and the other enclosing the entire list. This is perhap more easily seen in the following example, which creates two rewrite-db rules, one rewriting database dbA to dbB, and one rewriting database dbC to dbD:

```
CHANGE REPLICATION FILTER
  REPLICATE_REWRITE_DB = ((dbA, dbB), (dbC, dbD));
```

This statement leaves any existing replication filtering rules unchanged; to unset all filters of a given type, set the filter's value to an explicitly empty list, as shown in this example, which removes all existing REPLICATE_DO_DB and REPLICATE_IGNORE_DB rules:

```
CHANGE REPLICATION FILTER
    REPLICATE_DO_DB = (), REPLICATE_IGNORE_DB = ();
```

Setting a filter to empty in this way removes all existing rules, does not create any new ones, and does not restore any rules set at mysqld startup using --replicate-* options on the command line or in the configuration file.

Values employed with REPLICATE_WILD_DO_TABLE and REPLICATE_WILD_IGNORE_TABLE must be in the format *db_name.tbl_name*.

For more information, see Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”.

13.4.2.3 MASTER_POS_WAIT() Syntax

```
SELECT MASTER_POS_WAIT('master_log_file', master_log_pos [, timeout][, channel])
```

This is actually a function, not a statement. It is used to ensure that the slave has read and executed events up to a given position in the master's binary log. See Section 12.20, “Miscellaneous Functions”, for a full description.

13.4.2.4 RESET SLAVE Syntax

```
RESET SLAVE [ALL] [channel_option]

channel_option:
    FOR CHANNEL channel
```

RESET SLAVE makes the slave forget its replication position in the master's binary log. This statement is meant to be used for a clean start: It clears the master info and relay log info repositories, deletes all the relay log files, and starts a new relay log file. It also resets to 0 the replication delay specified with the MASTER_DELAY option to CHANGE MASTER TO. RESET SLAVE does not change the values of gtid_executed or gtid_purged. To use RESET SLAVE, the slave replication threads must be stopped, so on a running slave use STOP SLAVE before issuing RESET SLAVE.

Note

All relay log files are deleted, even if they have not been completely executed by the slave SQL thread. (This is a condition likely to exist on a replication slave if you have issued a STOP SLAVE statement or if the slave is highly loaded.)

The optional `FOR CHANNEL channel` clause enables you to choose which replication channel the statement applies to. If no clause is set and no extra channels exist, the statement applies to the default channel. Providing a `FOR CHANNEL channel` clause applies the `RESET SLAVE` statement to a specific replication channel. Combining a `FOR CHANNEL channel` clause with the `ALL` option deletes the specified channel. Issuing a `RESET SLAVE ALL` statement without a `FOR CHANNEL channel` clause when multiple replication channels exist deletes *all* replication channels and recreates only the default channel. See Section 18.2.3, "Replication Channels" for more information.

`RESET SLAVE` does not change any replication connection parameters such as master host, master port, master user, or master password, which are retained in memory. This means that `START SLAVE` can be issued without requiring a `CHANGE MASTER TO` statement following `RESET SLAVE`.

Connection parameters are reset by `RESET SLAVE ALL`. (`RESET SLAVE` followed by a restart of the slave `mysqld` also does this.)

In MySQL 8.0 `RESET SLAVE` causes an implicit commit of an ongoing transaction. See Section 13.3.3, "Statements That Cause an Implicit Commit".

If the slave SQL thread was in the middle of replicating temporary tables when it was stopped, and `RESET SLAVE` is issued, these replicated temporary tables are deleted on the slave.

Prior to MySQL 5.7.5, `RESET SLAVE` also had the effect of resetting both the heartbeat period (`Slave_heartbeat_period`) and `SSL_VERIFY_SERVER_CERT`. This issue is fixed in MySQL 5.7.5 and later. (Bug #18777899, Bug #18778485)

`RESET SLAVE ALL` clears the `IGNORE_SERVER_IDS` list set by `CHANGE MASTER TO`.

13.4.2.5 SET GLOBAL sql_slave_skip_counter Syntax

```
SET GLOBAL sql_slave_skip_counter = N
```

This statement skips the next *N* events from the master. This is useful for recovering from replication stops caused by a statement.

This statement is valid only when the slave threads are not running. Otherwise, it produces an error.

When using this statement, it is important to understand that the binary log is actually organized as a sequence of groups known as *event groups*. Each event group consists of a sequence of events.

- For transactional tables, an event group corresponds to a transaction.
- For nontransactional tables, an event group corresponds to a single SQL statement.

Note

A single transaction can contain changes to both transactional and nontransactional tables.

When you use `SET GLOBAL sql_slave_skip_counter` to skip events and the result is in the middle of a group, the slave continues to skip events until it reaches the end of the group. Execution then starts with the next event group.

13.4.2.6 START SLAVE Syntax

```
START SLAVE [thread_types] [until_option] [connection_options] [channel_option]

thread_types:
    [thread_type [, thread_type] ... ]

thread_type:
```

```
    IO_THREAD | SQL_THREAD

until_option:
    UNTIL {   {SQL_BEFORE_GTIDS | SQL_AFTER_GTIDS} = gtid_set
          |   MASTER_LOG_FILE = 'log_name', MASTER_LOG_POS = log_pos
          |   RELAY_LOG_FILE = 'log_name', RELAY_LOG_POS = log_pos
          |   SQL_AFTER_MTS_GAPS  }

connection_options:
    [USER='user_name'] [PASSWORD='user_pass'] [DEFAULT_AUTH='plugin_name'] [PLUGIN_DIR='plugin_dir']


channel_option:
    FOR CHANNEL channel

gtid_set:
    uuid_set [, uuid_set] ...
    | ''

uuid_set:
    uuid:interval[:interval]...

uuid:
    hhhhhhhh-hhhh-hhhh-hhhh-hhhhhhhhhhhh

h:
    [0-9,A-F]

interval:
    n[-n]

    (n >= 1)
```

`START SLAVE` with no *`thread_type`* options starts both of the slave threads. The I/O thread reads events from the master server and stores them in the relay log. The SQL thread reads events from the relay log and executes them. `START SLAVE` requires the `REPLICATION_SLAVE_ADMIN` or `SUPER` privilege.

If `START SLAVE` succeeds in starting the slave threads, it returns without any error. However, even in that case, it might be that the slave threads start and then later stop (for example, because they do not manage to connect to the master or read its binary log, or some other problem). `START SLAVE` does not warn you about this. You must check the slave's error log for error messages generated by the slave threads, or check that they are running satisfactorily with `SHOW SLAVE STATUS`.

`START SLAVE` causes an implicit commit of an ongoing transaction. See Section 13.3.3, “Statements That Cause an Implicit Commit”.

`gtid_next` must be set to `AUTOMATIC` before issuing this statement.

The optional `FOR CHANNEL` *`channel`* clause enables you to choose which replication channel the statement applies to. If no clause is set and no extra channels exist, the statement applies to the default channel and behaves the same as versions of MySQL prior to 5.7.6. Providing a `FOR CHANNEL` *`channel`* clause applies the `START SLAVE` statement to a specific replication channel. If a `START SLAVE` statement does not have a channel defined when using multiple channels, this statement starts the specified threads for all channels. This statement is disallowed for the `group_replication_recovery` channel. See Section 18.2.3, “Replication Channels” for more information.

MySQL 8.0 supports pluggable user-password authentication with `START SLAVE` with the `USER`, `PASSWORD`, `DEFAULT_AUTH` and `PLUGIN_DIR` options, as described in the following list:

- `USER`: User name. Cannot be set to an empty or null string, or left unset if `PASSWORD` is used.
- `PASSWORD`: Password.
- `DEFAULT_AUTH`: Name of plugin; default is MySQL native authentication.

- `PLUGIN_DIR`: Location of plugin.

You cannot use the `SQL_THREAD` option when specifying any of `USER`, `PASSWORD`, `DEFAULT_AUTH`, or `PLUGIN_DIR`, unless the `IO_THREAD` option is also provided.

See Section 6.3.10, “Pluggable Authentication”, for more information.

If an insecure connection is used with any these options, the server issues the warning `Sending passwords in plain text without SSL/TLS is extremely insecure`.

`START SLAVE ... UNTIL` supports two additional options for use with global transaction identifiers (GTIDs) (see Section 18.1.3, “Replication with Global Transaction Identifiers”). Each of these takes a set of one or more global transaction identifiers *`gtid_set`* as an argument (see GTID Sets, for more information).

When no *`thread_type`* is specified, `START SLAVE UNTIL SQL_BEFORE_GTIDS` causes the slave SQL thread to process transactions until it has reached the *first* transaction whose GTID is listed in the *`gtid_set`*. `START SLAVE UNTIL SQL_AFTER_GTIDS` causes the slave threads to process all transactions until the *`last`* transaction in the *`gtid_set`* has been processed by both threads. In other words, `START SLAVE UNTIL SQL_BEFORE_GTIDS` causes the slave SQL thread to process all transactions occurring before the first GTID in the *`gtid_set`* is reached, and `START SLAVE UNTIL SQL_AFTER_GTIDS` causes the slave threads to handle all transactions, including those whose GTIDs are found in *`gtid_set`*, until each has encountered a transaction whose GTID is not part of the set. `SQL_BEFORE_GTIDS` and `SQL_AFTER_GTIDS` each support the `SQL_THREAD` and `IO_THREAD` options, although using `IO_THREAD` with them currently has no effect.

For example, `START SLAVE SQL_THREAD UNTIL SQL_BEFORE_GTIDS = 3E11FA47-71CA-11E1-9E33-C80AA9429562:11-56` causes the slave SQL thread to process all transactions originating from the master whose `server_uuid` is `3E11FA47-71CA-11E1-9E33-C80AA9429562` until it encounters the transaction having sequence number 11; it then stops without processing this transaction. In other words, all transactions up to and including the transaction with sequence number 10 are processed. Executing `START SLAVE SQL_THREAD UNTIL SQL_AFTER_GTIDS = 3E11FA47-71CA-11E1-9E33-C80AA9429562:11-56`, on the other hand, would cause the slave SQL thread to obtain all transactions just mentioned from the master, including all of the transactions having the sequence numbers 11 through 56, and then to stop without processing any additional transactions; that is, the transaction having sequence number 56 would be the last transaction fetched by the slave SQL thread.

When using a multi-threaded slave, there is a chance of gaps in the sequence of transactions that have been executed from the relay log in the following cases:

- killing the coordinator thread
- after an error occurs in the worker threads
- `mysqld` shuts down unexpectedly

Use the `START SLAVE UNTIL SQL_AFTER_MTS_GAPS` statement to cause a multi-threaded slave's worker threads to only run until no more gaps are found in the relay log, and then to stop. This statement can take an `SQL_THREAD` option, but the effects of the statement remain unchanged. It has no effect on the slave I/O thread (and cannot be used with the `IO_THREAD` option).

Issuing `START SLAVE` on a multi-threaded slave with gaps in the sequence of transactions executed from the relay log generates a warning. In such a situation, the solution is to use `START SLAVE UNTIL SQL_AFTER_MTS_GAPS`, then issue `RESET SLAVE` to remove any remaining relay logs. See Section 18.4.1.34, “Replication and Transaction Inconsistencies” for more information.

To change a failed multi-threaded slave to single-threaded mode, you can issue the following series of statements, in the order shown:

```
START SLAVE UNTIL SQL_AFTER_MTS_GAPS;

SET @@GLOBAL.slave_parallel_workers = 0;

START SLAVE SQL_THREAD;
```

Note

It is possible to view the entire text of a running `START SLAVE ...` statement, including any `USER` or `PASSWORD` values used, in the output of `SHOW PROCESSLIST`. This is also true for the text of a running `CHANGE MASTER TO` statement, including any values it employs for `MASTER_USER` or `MASTER_PASSWORD`.

`START SLAVE` sends an acknowledgment to the user after both the I/O thread and the SQL thread have started. However, the I/O thread may not yet have connected. For this reason, a successful `START SLAVE` causes `SHOW SLAVE STATUS` to show `Slave_SQL_Running=Yes`, but this does not guarantee that `Slave_IO_Running=Yes` (because `Slave_IO_Running=Yes` only if the I/O thread is running *and connected*). For more information, see Section 13.7.5.34, “SHOW SLAVE STATUS Syntax”, and Section 18.1.7.1, “Checking Replication Status”.

You can add `IO_THREAD` and `SQL_THREAD` options to the statement to name which of the threads to start. The `SQL_THREAD` option is disallowed when specifying any of `USER`, `PASSWORD`, `DEFAULT_AUTH`, or `PLUGIN_DIR`, unless the `IO_THREAD` option is also provided.

An `UNTIL` clause (*`until_option`*, in the preceding grammar) may be added to specify that the slave should start and run until the SQL thread reaches a given point in the master binary log, specified by the `MASTER_LOG_POS` and MASTER_LOG_FILE options, or a given point in the slave relay log, indicated with the `RELAY_LOG_POS` and `RELAY_LOG_FILE` options. When the SQL thread reaches the point specified, it stops. If the `SQL_THREAD` option is specified in the statement, it starts only the SQL thread. Otherwise, it starts both slave threads. If the SQL thread is running, the `UNTIL` clause is ignored and a warning is issued. You cannot use an `UNTIL` clause with the `IO_THREAD` option.

It is also possible with `START SLAVE UNTIL` to specify a stop point relative to a given GTID or set of GTIDs using one of the options `SQL_BEFORE_GTIDS` or `SQL_AFTER_GTIDS`, as explained previously in this section. When using one of these options, you can specify `SQL_THREAD`, `IO_THREAD`, both of these, or neither of them. If you specify only `SQL_THREAD`, then only the slave SQL thread is affected by the statement; if only `IO_THREAD` is used, then only the slave I/O is affected. If both `SQL_THREAD` and `IO_THREAD` are used, or if neither of them is used, then both the SQL and I/O threads are affected by the statement.

The `UNTIL` clause is not supported for multi-threaded slaves except when also using `SQL_AFTER_MTS_GAPS`.

For an `UNTIL` clause, you must specify any one of the following:

- *Both* a log file name and a position in that file
- *Either* of `SQL_BEFORE_GTIDS` or `SQL_AFTER_GTIDS`
- `SQL_AFTER_MTS_GAPS`

Do not mix master and relay log options. Do not mix log file options with GTID options.

Any `UNTIL` condition is reset by a subsequent `STOP SLAVE` statement, a `START SLAVE` statement that includes no `UNTIL` clause, or a server restart.

When specifying a log file and position, you can use the `IO_THREAD` option with `START SLAVE ... UNTIL` even though only the SQL thread is affected by this statement. The `IO_THREAD` option is ignored in such cases. The preceding restriction does not apply when using one of the GTID options (`SQL_BEFORE_GTIDS` and `SQL_AFTER_GTIDS`); the GTID options support both `SQL_THREAD` and `IO_THREAD`, as explained previously in this section.

The UNTIL clause can be useful for debugging replication, or to cause replication to proceed until just before the point where you want to avoid having the slave replicate an event. For example, if an unwise DROP TABLE statement was executed on the master, you can use UNTIL to tell the slave to execute up to that point but no farther. To find what the event is, use mysqlbinlog with the master binary log or slave relay log, or by using a SHOW BINLOG EVENTS statement.

If you are using UNTIL to have the slave process replicated queries in sections, it is recommended that you start the slave with the --skip-slave-start option to prevent the SQL thread from running when the slave server starts. It is probably best to use this option in an option file rather than on the command line, so that an unexpected server restart does not cause it to be forgotten.

The SHOW SLAVE STATUS statement includes output fields that display the current values of the UNTIL condition.

In very old versions of MySQL (before 4.0.5), this statement was called SLAVE START. In MySQL 8.0, that syntax produces an error.

13.4.2.7 STOP SLAVE Syntax

```
STOP SLAVE [thread_types]

thread_types:
    [thread_type [, thread_type] ... ]

thread_type: IO_THREAD | SQL_THREAD

channel_option:
    FOR CHANNEL channel
```

Stops the slave threads. STOP SLAVE requires the REPLICATION_SLAVE_ADMIN or SUPER privilege. Recommended best practice is to execute STOP SLAVE on the slave before stopping the slave server (see Section 5.1.12, “The Server Shutdown Process”, for more information).

When using the row-based logging format: You should execute STOP SLAVE or STOP SLAVE SQL_THREAD on the slave prior to shutting down the slave server if you are replicating any tables that use a nontransactional storage engine (see the *Note* later in this section).

Like START SLAVE, this statement may be used with the IO_THREAD and SQL_THREAD options to name the thread or threads to be stopped.

In MySQL 8.0, STOP SLAVE causes an implicit commit of an ongoing transaction. See Section 13.3.3, “Statements That Cause an Implicit Commit”.

gtid_next must be set to AUTOMATIC before issuing this statement.

You can control how long STOP SLAVE waits before timing out by setting the rpl_stop_slave_timeout system variable. This can be used to avoid deadlocks between STOP SLAVE and other slave SQL statements using different client connections to the slave.

Some CHANGE MASTER TO statements are allowed while the slave is running, depending on the states of the slave SQL and I/O threads. However, using STOP SLAVE prior to executing CHANGE MASTER TO in such cases is still supported. See Section 13.4.2.1, “CHANGE MASTER TO Syntax”, and Section 18.3.8, “Switching Masters During Failover”, for more information.

The optional FOR CHANNEL *channel* clause enables you to choose which replication channel the statement applies to. If no clause is set and no extra channels exist, the statement applies to the default channel and behaves the same as versions of MySQL prior to 5.7. Providing a FOR CHANNEL *channel* clause applies the STOP SLAVE statement to a specific replication channel. If a STOP SLAVE statement does not have a channel defined when using multiple channels, this statement stops the specified threads for all channels. This statement cannot be used with the

`group_replication_recovery` channel. See Section 18.2.3, “Replication Channels” for more information.

When using statement-based replication: changing the master while it has open temporary tables is potentially unsafe. This is one of the reasons why statement-based replication of temporary tables is not recommended. You can find out whether there are any temporary tables on the slave by checking the value of `Slave_open_temp_tables`; when using statement-based replication, this value should be 0 before executing `CHANGE MASTER TO`. If there are any temporary tables open on the slave, issuing a `CHANGE MASTER TO` statement after issuing a `STOP SLAVE` causes an **`ER_WARN_OPEN_TEMP_TABLES_MUST_BE_ZERO`** warning.

When using a multi-threaded slave (`slave_parallel_workers` is a nonzero value), any gaps in the sequence of transactions executed from the relay log are closed as part of stopping the worker threads. If the slave is stopped unexpectedly (for example due to an error in a worker thread, or another thread issuing `KILL`) while a `STOP SLAVE` statement is executing, the sequence of executed transactions from the relay log may become inconsistent. See Section 18.4.1.34, “Replication and Transaction Inconsistencies” for more information.

In MySQL 8.0, `STOP SLAVE` waits until the current replication transaction affecting one or more nontransactional tables has finished executing (if there is any such transaction), or until you issue a `KILL QUERY` or `KILL CONNECTION` statement. (Bug #319, Bug #38205)

13.4.3 SQL Statements for Controlling Group Replication

This section provides information about the statements used for controlling group replication.

13.4.3.1 START GROUP_REPLICATION Syntax

```
START GROUP_REPLICATION
```

Starts group replication. This statement requires the `GROUP_REPLICATION_ADMIN` or `SUPER` privilege.

13.4.3.2 STOP GROUP_REPLICATION Syntax

```
STOP GROUP_REPLICATION
```

Stops group replication. This statement requires the `GROUP_REPLICATION_ADMIN` or `SUPER` privilege.

13.5 Prepared SQL Statement Syntax

MySQL 8.0 provides support for server-side prepared statements. This support takes advantage of the efficient client/server binary protocol. Using prepared statements with placeholders for parameter values has the following benefits:

- Less overhead for parsing the statement each time it is executed. Typically, database applications process large volumes of almost-identical statements, with only changes to literal or variable values in clauses such as `WHERE` for queries and deletes, `SET` for updates, and `VALUES` for inserts.
- Protection against SQL injection attacks. The parameter values can contain unescaped SQL quote and delimiter characters.

Prepared Statements in Application Programs

You can use server-side prepared statements through client programming interfaces, including the MySQL C API client library or MySQL Connector/C for C programs, MySQL Connector/J for Java programs, and MySQL Connector/Net for programs using .NET technologies. For example, the C API provides a set of function calls that make up its prepared statement API. See Section 27.7.8, “C API Prepared Statements”. Other language interfaces can provide support for prepared statements that use

the binary protocol by linking in the C client library, one example being the `mysqli` extension, available in PHP 5.0 and later.

Prepared Statements in SQL Scripts

An alternative SQL interface to prepared statements is available. This interface is not as efficient as using the binary protocol through a prepared statement API, but requires no programming because it is available directly at the SQL level:

- You can use it when no programming interface is available to you.
- You can use it from any program that can send SQL statements to the server to be executed, such as the `mysql` client program.
- You can use it even if the client is using an old version of the client library, as long as you connect to a server running MySQL 4.1 or higher.

SQL syntax for prepared statements is intended to be used for situations such as these:

- To test how prepared statements work in your application before coding it.
- To use prepared statements when you do not have access to a programming API that supports them.
- To interactively troubleshoot application issues with prepared statements.
- To create a test case that reproduces a problem with prepared statements, so that you can file a bug report.

PREPARE, EXECUTE, and DEALLOCATE PREPARE Statements

SQL syntax for prepared statements is based on three SQL statements:

- `PREPARE` prepares a statement for execution (see Section 13.5.1, “PREPARE Syntax”).
- `EXECUTE` executes a prepared statement (see Section 13.5.2, “EXECUTE Syntax”).
- `DEALLOCATE PREPARE` releases a prepared statement (see Section 13.5.3, “DEALLOCATE PREPARE Syntax”).

The following examples show two equivalent ways of preparing a statement that computes the hypotenuse of a triangle given the lengths of the two sides.

The first example shows how to create a prepared statement by using a string literal to supply the text of the statement:

```
mysql> PREPARE stmt1 FROM 'SELECT SQRT(POW(?,2) + POW(?,2)) AS hypotenuse';
mysql> SET @a = 3;
mysql> SET @b = 4;
mysql> EXECUTE stmt1 USING @a, @b;
+------------+
| hypotenuse |
+------------+
|          5 |
+------------+
mysql> DEALLOCATE PREPARE stmt1;
```

The second example is similar, but supplies the text of the statement as a user variable:

```
mysql> SET @s = 'SELECT SQRT(POW(?,2) + POW(?,2)) AS hypotenuse';
mysql> PREPARE stmt2 FROM @s;
```

```
mysql> SET @a = 6;
mysql> SET @b = 8;
mysql> EXECUTE stmt2 USING @a, @b;
+------------+
| hypotenuse |
+------------+
|         10 |
+------------+
mysql> DEALLOCATE PREPARE stmt2;
```

Here is an additional example that demonstrates how to choose the table on which to perform a query at runtime, by storing the name of the table as a user variable:

```
mysql> USE test;
mysql> CREATE TABLE t1 (a INT NOT NULL);
mysql> INSERT INTO t1 VALUES (4), (8), (11), (32), (80);

mysql> SET @table = 't1';
mysql> SET @s = CONCAT('SELECT * FROM ', @table);

mysql> PREPARE stmt3 FROM @s;
mysql> EXECUTE stmt3;
+----+
| a  |
+----+
|  4 |
|  8 |
| 11 |
| 32 |
| 80 |
+----+

mysql> DEALLOCATE PREPARE stmt3;
```

A prepared statement is specific to the session in which it was created. If you terminate a session without deallocating a previously prepared statement, the server deallocates it automatically.

A prepared statement is also global to the session. If you create a prepared statement within a stored routine, it is not deallocated when the stored routine ends.

To guard against too many prepared statements being created simultaneously, set the `max_prepared_stmt_count` system variable. To prevent the use of prepared statements, set the value to 0.

SQL Syntax Allowed in Prepared Statements

The following SQL statements can be used as prepared statements:

```
ALTER TABLE
ALTER USER
ANALYZE TABLE
CACHE INDEX
CALL
CHANGE MASTER
CHECKSUM {TABLE | TABLES}
COMMIT
{CREATE | DROP} INDEX
{CREATE | RENAME | DROP} DATABASE
{CREATE | DROP} TABLE
{CREATE | RENAME | DROP} USER
{CREATE | DROP} VIEW
DELETE
DO
FLUSH {TABLE | TABLES | TABLES WITH READ LOCK | HOSTS | PRIVILEGES
  | LOGS | STATUS | MASTER | SLAVE | DES_KEY_FILE | USER_RESOURCES}
GRANT
INSERT
```

```
INSTALL PLUGIN
KILL
LOAD INDEX INTO CACHE
OPTIMIZE TABLE
RENAME TABLE
REPAIR TABLE
REPLACE
RESET {MASTER | SLAVE | QUERY CACHE}
REVOKE
SELECT
SET
SHOW {WARNINGS | ERRORS}
SHOW BINLOG EVENTS
SHOW CREATE {PROCEDURE | FUNCTION | EVENT | TABLE | VIEW}
SHOW {MASTER | BINARY} LOGS
SHOW {MASTER | SLAVE} STATUS
SLAVE {START | STOP}
TRUNCATE TABLE
UNINSTALL PLUGIN
UPDATE
```

For compliance with the SQL standard, which states that diagnostics statements are not preparable, MySQL does not support the following as prepared statements:

- `SHOW WARNINGS`, `SHOW COUNT(*) WARNINGS`
- `SHOW ERRORS`, `SHOW COUNT(*) ERRORS`
- Statements containing any reference to the `warning_count` or `error_count` system variable.

Other statements are not supported in MySQL 8.0.

Generally, statements not permitted in SQL prepared statements are also not permitted in stored programs. Exceptions are noted in Section C.1, “Restrictions on Stored Programs”.

Metadata changes to tables or views referred to by prepared statements are detected and cause automatic repreparation of the statement when it is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

Placeholders can be used for the arguments of the `LIMIT` clause when using prepared statements. See Section 13.2.10, “SELECT Syntax”.

In prepared `CALL` statements used with `PREPARE` and `EXECUTE`, placeholder support for `OUT` and `INOUT` parameters is available beginning with MySQL 8.0. See Section 13.2.1, “CALL Syntax”, for an example and a workaround for earlier versions. Placeholders can be used for `IN` parameters regardless of version.

SQL syntax for prepared statements cannot be used in nested fashion. That is, a statement passed to `PREPARE` cannot itself be a `PREPARE`, `EXECUTE`, or `DEALLOCATE PREPARE` statement.

SQL syntax for prepared statements is distinct from using prepared statement API calls. For example, you cannot use the `mysql_stmt_prepare()` C API function to prepare a `PREPARE`, `EXECUTE`, or `DEALLOCATE PREPARE` statement.

SQL syntax for prepared statements can be used within stored procedures, but not in stored functions or triggers. However, a cursor cannot be used for a dynamic statement that is prepared and executed with `PREPARE` and `EXECUTE`. The statement for a cursor is checked at cursor creation time, so the statement cannot be dynamic.

SQL syntax for prepared statements does not support multi-statements (that is, multiple statements within a single string separated by `;` characters).

Prepared statements use the query cache under the conditions described in Section 8.10.3.1, “How the Query Cache Operates”.

To write C programs that use the CALL SQL statement to execute stored procedures that contain prepared statements, the CLIENT_MULTI_RESULTS flag must be enabled. This is because each CALL returns a result to indicate the call status, in addition to any result sets that might be returned by statements executed within the procedure.

CLIENT_MULTI_RESULTS can be enabled when you call mysql_real_connect(), either explicitly by passing the CLIENT_MULTI_RESULTS flag itself, or implicitly by passing CLIENT_MULTI_STATEMENTS (which also enables CLIENT_MULTI_RESULTS). For additional information, see Section 13.2.1, “CALL Syntax”.

13.5.1 PREPARE Syntax

```
PREPARE stmt_name FROM preparable_stmt
```

The PREPARE statement prepares a SQL statement and assigns it a name, *stmt_name*, by which to refer to the statement later. The prepared statement is executed with EXECUTE and released with DEALLOCATE PREPARE. For examples, see Section 13.5, “Prepared SQL Statement Syntax”.

Statement names are not case sensitive. *preparable_stmt* is either a string literal or a user variable that contains the text of the SQL statement. The text must represent a single statement, not multiple statements. Within the statement, ? characters can be used as parameter markers to indicate where data values are to be bound to the query later when you execute it. The ? characters should not be enclosed within quotation marks, even if you intend to bind them to string values. Parameter markers can be used only where data values should appear, not for SQL keywords, identifiers, and so forth.

If a prepared statement with the given name already exists, it is deallocated implicitly before the new statement is prepared. This means that if the new statement contains an error and cannot be prepared, an error is returned and no statement with the given name exists.

The scope of a prepared statement is the session within which it is created, which as several implications:

- A prepared statement created in one session is not available to other sessions.
- When a session ends, whether normally or abnormally, its prepared statements no longer exist. If auto-reconnect is enabled, the client is not notified that the connection was lost. For this reason, clients may wish to disable auto-reconnect. See Section 27.7.19, “Controlling Automatic Reconnection Behavior”.
- A prepared statement created within a stored program continues to exist after the program finishes executing and can be executed outside the program later.
- A statement prepared in stored program context cannot refer to stored procedure or function parameters or local variables because they go out of scope when the program ends and would be unavailable were the statement to be executed later outside the program. As a workaround, refer instead to user-defined variables, which also have session scope; see Section 9.4, “User-Defined Variables”.

13.5.2 EXECUTE Syntax

```
EXECUTE stmt_name
    [USING @var_name [, @var_name] ...]
```

After preparing a statement with PREPARE, you execute it with an EXECUTE statement that refers to the prepared statement name. If the prepared statement contains any parameter markers, you must supply a USING clause that lists user variables containing the values to be bound to the parameters. Parameter values can be supplied only by user variables, and the USING clause must name exactly as many variables as the number of parameter markers in the statement.

You can execute a given prepared statement multiple times, passing different variables to it or setting the variables to different values before each execution.

For examples, see Section 13.5, “Prepared SQL Statement Syntax”.

13.5.3 DEALLOCATE PREPARE Syntax

```
{DEALLOCATE | DROP} PREPARE stmt_name
```

To deallocate a prepared statement produced with `PREPARE`, use a `DEALLOCATE PREPARE` statement that refers to the prepared statement name. Attempting to execute a prepared statement after deallocating it results in an error. If too many prepared statements are created and not deallocated by either the `DEALLOCATE PREPARE` statement or the end of the session, you might encounter the upper limit enforced by the `max_prepared_stmt_count` system variable.

For examples, see Section 13.5, “Prepared SQL Statement Syntax”.

13.6 Compound-Statement Syntax

This section describes the syntax for the `BEGIN ... END` compound statement and other statements that can be used in the body of stored programs: Stored procedures and functions, triggers, and events. These objects are defined in terms of SQL code that is stored on the server for later invocation (see Chapter 23, *Stored Programs and Views*).

A compound statement is a block that can contain other blocks; declarations for variables, condition handlers, and cursors; and flow control constructs such as loops and conditional tests.

13.6.1 BEGIN ... END Compound-Statement Syntax

```
[begin_label:] BEGIN
    [statement_list]
END [end_label]
```

`BEGIN ... END` syntax is used for writing compound statements, which can appear within stored programs (stored procedures and functions, triggers, and events). A compound statement can contain multiple statements, enclosed by the `BEGIN` and `END` keywords. *`statement_list`* represents a list of one or more statements, each terminated by a semicolon (`;`) statement delimiter. The *`statement_list`* itself is optional, so the empty compound statement (`BEGIN END`) is legal.

`BEGIN ... END` blocks can be nested.

Use of multiple statements requires that a client is able to send statement strings containing the `;` statement delimiter. In the `mysql` command-line client, this is handled with the `delimiter` command. Changing the `;` end-of-statement delimiter (for example, to `//`) permit `;` to be used in a program body. For an example, see Section 23.1, “Defining Stored Programs”.

A `BEGIN ... END` block can be labeled. See Section 13.6.2, “Statement Label Syntax”.

The optional `[NOT] ATOMIC` clause is not supported. This means that no transactional savepoint is set at the start of the instruction block and the `BEGIN` clause used in this context has no effect on the current transaction.

Note

Within all stored programs, the parser treats `BEGIN [WORK]` as the beginning of a `BEGIN ... END` block. To begin a transaction in this context, use `START TRANSACTION` instead.

13.6.2 Statement Label Syntax

```
[begin_label:] BEGIN
    [statement_list]
END [end_label]

[begin_label:] LOOP
    statement_list
END LOOP [end_label]

[begin_label:] REPEAT
    statement_list
UNTIL search_condition
END REPEAT [end_label]

[begin_label:] WHILE search_condition DO
    statement_list
END WHILE [end_label]
```

Labels are permitted for `BEGIN ... END` blocks and for the `LOOP`, `REPEAT`, and `WHILE` statements. Label use for those statements follows these rules:

- *`begin_label`* must be followed by a colon.
- *`begin_label`* can be given without *`end_label`*. If *`end_label`* is present, it must be the same as *`begin_label`*.
- *`end_label`* cannot be given without *`begin_label`*.
- Labels at the same nesting level must be distinct.
- Labels can be up to 16 characters long.

To refer to a label within the labeled construct, use an `ITERATE` or `LEAVE` statement. The following example uses those statements to continue iterating or terminate the loop:

```
CREATE PROCEDURE doiterate(p1 INT)
BEGIN
  label1: LOOP
    SET p1 = p1 + 1;
    IF p1 < 10 THEN ITERATE label1; END IF;
    LEAVE label1;
  END LOOP label1;
END;
```

The scope of a block label does not include the code for handlers declared within the block. For details, see Section 13.6.7.2, “DECLARE ... HANDLER Syntax”.

13.6.3 DECLARE Syntax

The `DECLARE` statement is used to define various items local to a program:

- Local variables. See Section 13.6.4, “Variables in Stored Programs”.
- Conditions and handlers. See Section 13.6.7, “Condition Handling”.
- Cursors. See Section 13.6.6, “Cursors”.

`DECLARE` is permitted only inside a `BEGIN ... END` compound statement and must be at its start, before any other statements.

Declarations must follow a certain order. Cursor declarations must appear before handler declarations. Variable and condition declarations must appear before cursor or handler declarations.

13.6.4 Variables in Stored Programs

System variables and user-defined variables can be used in stored programs, just as they can be used outside stored-program context. In addition, stored programs can use `DECLARE` to define local variables, and stored routines (procedures and functions) can be declared to take parameters that communicate values between the routine and its caller.

- To declare local variables, use the `DECLARE` statement, as described in Section 13.6.4.1, “Local Variable DECLARE Syntax”.
- Variables can be set directly with the `SET` statement. See Section 13.7.4.1, “SET Syntax for Variable Assignment”.
- Results from queries can be retrieved into local variables using `SELECT ... INTO var_list` or by opening a cursor and using `FETCH ... INTO var_list`. See Section 13.2.10.1, “SELECT ... INTO Syntax”, and Section 13.6.6, “Cursors”.

For information about the scope of local variables and how MySQL resolves ambiguous names, see Section 13.6.4.2, “Local Variable Scope and Resolution”.

It is not permitted to assign the value `DEFAULT` to stored procedure or function parameters or stored program local variables (for example with a `SET var_name = DEFAULT` statement). In MySQL 8.0, this results in a syntax error.

13.6.4.1 Local Variable DECLARE Syntax

```
DECLARE var_name [, var_name] ... type [DEFAULT value]
```

This statement declares local variables within stored programs. To provide a default value for a variable, include a `DEFAULT` clause. The value can be specified as an expression; it need not be a constant. If the `DEFAULT` clause is missing, the initial value is `NULL`.

Local variables are treated like stored routine parameters with respect to data type and overflow checking. See Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”.

Variable declarations must appear before cursor or handler declarations.

Local variable names are not case sensitive. Permissible characters and quoting rules are the same as for other identifiers, as described in Section 9.2, “Schema Object Names”.

The scope of a local variable is the `BEGIN ... END` block within which it is declared. The variable can be referred to in blocks nested within the declaring block, except those blocks that declare a variable with the same name.

For examples of variable declarations, see Section 13.6.4.2, “Local Variable Scope and Resolution”.

13.6.4.2 Local Variable Scope and Resolution

The scope of a local variable is the `BEGIN ... END` block within which it is declared. The variable can be referred to in blocks nested within the declaring block, except those blocks that declare a variable with the same name.

Because local variables are in scope only during stored program execution, references to them are not permitted in prepared statements created within a stored program. Prepared statement scope is the current session, not the stored program, so the statement could be executed after the program ends, at which point the variables would no longer be in scope. For example, `SELECT ... INTO local_var` cannot be used as a prepared statement. This restriction also applies to stored procedure and function parameters. See Section 13.5.1, “PREPARE Syntax”.

A local variable should not have the same name as a table column. If an SQL statement, such as a `SELECT ... INTO` statement, contains a reference to a column and a declared local variable with the same name, MySQL currently interprets the reference as the name of a variable. Consider the following procedure definition:

```
CREATE PROCEDURE sp1 (x VARCHAR(5))
BEGIN
  DECLARE xname VARCHAR(5) DEFAULT 'bob';
  DECLARE newname VARCHAR(5);
  DECLARE xid INT;

  SELECT xname, id INTO newname, xid
    FROM table1 WHERE xname = xname;
  SELECT newname;
END;
```

MySQL interprets `xname` in the `SELECT` statement as a reference to the `xname` *variable* rather than the `xname` *column*. Consequently, when the procedure `sp1()`is called, the `newname` variable returns the value `'bob'` regardless of the value of the `table1.xname` column.

Similarly, the cursor definition in the following procedure contains a `SELECT` statement that refers to `xname`. MySQL interprets this as a reference to the variable of that name rather than a column reference.

```
CREATE PROCEDURE sp2 (x VARCHAR(5))
BEGIN
  DECLARE xname VARCHAR(5) DEFAULT 'bob';
  DECLARE newname VARCHAR(5);
  DECLARE xid INT;
  DECLARE done TINYINT DEFAULT 0;
  DECLARE cur1 CURSOR FOR SELECT xname, id FROM table1;
  DECLARE CONTINUE HANDLER FOR NOT FOUND SET done = 1;

  OPEN cur1;
  read_loop: LOOP
    FETCH FROM cur1 INTO newname, xid;
    IF done THEN LEAVE read_loop; END IF;
    SELECT newname;
  END LOOP;
  CLOSE cur1;
END;
```

See also Section C.1, “Restrictions on Stored Programs”.

13.6.5 Flow Control Statements

MySQL supports the `IF`, `CASE`, `ITERATE`, `LEAVE LOOP`, `WHILE`, and `REPEAT` constructs for flow control within stored programs. It also supports `RETURN` within stored functions.

Many of these constructs contain other statements, as indicated by the grammar specifications in the following sections. Such constructs may be nested. For example, an `IF` statement might contain a `WHILE` loop, which itself contains a `CASE` statement.

MySQL does not support `FOR` loops.

13.6.5.1 CASE Syntax

```
CASE case_value
    WHEN when_value THEN statement_list
    [WHEN when_value THEN statement_list] ...
    [ELSE statement_list]
END CASE
```

Or:

```
CASE
    WHEN search_condition THEN statement_list
    [WHEN search_condition THEN statement_list] ...
    [ELSE statement_list]
END CASE
```

The CASE statement for stored programs implements a complex conditional construct.

Note

There is also a CASE *expression*, which differs from the CASE *statement* described here. See Section 12.4, “Control Flow Functions”. The CASE statement cannot have an ELSE NULL clause, and it is terminated with END CASE instead of END.

For the first syntax, *case_value* is an expression. This value is compared to the *when_value* expression in each WHEN clause until one of them is equal. When an equal *when_value* is found, the corresponding THEN clause *statement_list* executes. If no *when_value* is equal, the ELSE clause *statement_list* executes, if there is one.

This syntax cannot be used to test for equality with NULL because NULL = NULL is false. See Section 3.3.4.6, “Working with NULL Values”.

For the second syntax, each WHEN clause *search_condition* expression is evaluated until one is true, at which point its corresponding THEN clause *statement_list* executes. If no *search_condition* is equal, the ELSE clause *statement_list* executes, if there is one.

If no *when_value* or *search_condition* matches the value tested and the CASE statement contains no ELSE clause, a Case not found for CASE statement error results.

Each *statement_list* consists of one or more SQL statements; an empty *statement_list* is not permitted.

To handle situations where no value is matched by any WHEN clause, use an ELSE containing an empty BEGIN ... END block, as shown in this example. (The indentation used here in the ELSE clause is for purposes of clarity only, and is not otherwise significant.)

```
DELIMITER |

CREATE PROCEDURE p()
  BEGIN
    DECLARE v INT DEFAULT 1;

    CASE v
      WHEN 2 THEN SELECT v;
      WHEN 3 THEN SELECT 0;
      ELSE
        BEGIN
        END;
    END CASE;
  END;
  |
```

13.6.5.2 IF Syntax

```
IF search_condition THEN statement_list
    [ELSEIF search_condition THEN statement_list] ...
    [ELSE statement_list]
END IF
```

The IF statement for stored programs implements a basic conditional construct.

Note

There is also an `IF()` *function*, which differs from the `IF` *statement* described here. See Section 12.4, “Control Flow Functions”. The `IF` statement can have `THEN`, `ELSE`, and `ELSEIF` clauses, and it is terminated with `END IF`.

If the *`search_condition`* evaluates to true, the corresponding `THEN` or `ELSEIF` clause *`statement_list`* executes. If no *`search_condition`* matches, the `ELSE` clause *`statement_list`* executes.

Each *`statement_list`* consists of one or more SQL statements; an empty *`statement_list`* is not permitted.

An `IF ... END IF` block, like all other flow-control blocks used within stored programs, must be terminated with a semicolon, as shown in this example:

```
DELIMITER //

CREATE FUNCTION SimpleCompare(n INT, m INT)
  RETURNS VARCHAR(20)

  BEGIN
    DECLARE s VARCHAR(20);

    IF n > m THEN SET s = '>';
    ELSEIF n = m THEN SET s = '=';
    ELSE SET s = '<';
    END IF;

    SET s = CONCAT(n, ' ', s, ' ', m);

    RETURN s;
  END //

DELIMITER ;
```

As with other flow-control constructs, `IF ... END IF` blocks may be nested within other flow-control constructs, including other `IF` statements. Each `IF` must be terminated by its own `END IF` followed by a semicolon. You can use indentation to make nested flow-control blocks more easily readable by humans (although this is not required by MySQL), as shown here:

```
DELIMITER //

CREATE FUNCTION VerboseCompare (n INT, m INT)
  RETURNS VARCHAR(50)

  BEGIN
    DECLARE s VARCHAR(50);

    IF n = m THEN SET s = 'equals';
    ELSE
      IF n > m THEN SET s = 'greater';
      ELSE SET s = 'less';
      END IF;

      SET s = CONCAT('is ', s, ' than');
    END IF;

    SET s = CONCAT(n, ' ', s, ' ', m, '.');

    RETURN s;
  END //

DELIMITER ;
```

In this example, the inner `IF` is evaluated only if `n` is not equal to `m`.

13.6.5.3 ITERATE Syntax

```
ITERATE label
```

`ITERATE` can appear only within `LOOP`, `REPEAT`, and `WHILE` statements. `ITERATE` means “start the loop again.”

For an example, see Section 13.6.5.5, “LOOP Syntax”.

13.6.5.4 LEAVE Syntax

```
LEAVE label
```

This statement is used to exit the flow control construct that has the given label. If the label is for the outermost stored program block, `LEAVE` exits the program.

`LEAVE` can be used within `BEGIN ... END` or loop constructs (`LOOP`, `REPEAT`, `WHILE`).

For an example, see Section 13.6.5.5, “LOOP Syntax”.

13.6.5.5 LOOP Syntax

```
[begin_label:] LOOP
    statement_list
END LOOP [end_label]
```

`LOOP` implements a simple loop construct, enabling repeated execution of the statement list, which consists of one or more statements, each terminated by a semicolon (`;`) statement delimiter. The statements within the loop are repeated until the loop is terminated. Usually, this is accomplished with a `LEAVE` statement. Within a stored function, `RETURN` can also be used, which exits the function entirely.

Neglecting to include a loop-termination statement results in an infinite loop.

A `LOOP` statement can be labeled. For the rules regarding label use, see Section 13.6.2, “Statement Label Syntax”.

Example:

```
CREATE PROCEDURE doiterate(p1 INT)
BEGIN
  label1: LOOP
    SET p1 = p1 + 1;
    IF p1 < 10 THEN
      ITERATE label1;
    END IF;
    LEAVE label1;
  END LOOP label1;
  SET @x = p1;
END;
```

13.6.5.6 REPEAT Syntax

```
[begin_label:] REPEAT
    statement_list
UNTIL search_condition
END REPEAT [end_label]
```

The statement list within a `REPEAT` statement is repeated until the *`search_condition`* expression is true. Thus, a `REPEAT` always enters the loop at least once. *`statement_list`* consists of one or more statements, each terminated by a semicolon (`;`) statement delimiter.

A REPEAT statement can be labeled. For the rules regarding label use, see Section 13.6.2, “Statement Label Syntax”.

Example:

```
mysql> delimiter //

mysql> CREATE PROCEDURE dorepeat(p1 INT)
    -> BEGIN
    ->   SET @x = 0;
    ->   REPEAT
    ->     SET @x = @x + 1;
    ->   UNTIL @x > p1 END REPEAT;
    -> END
    -> //
Query OK, 0 rows affected (0.00 sec)

mysql> CALL dorepeat(1000)//
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @x//
+------+
| @x   |
+------+
| 1001 |
+------+
1 row in set (0.00 sec)
```

13.6.5.7 RETURN Syntax

```
RETURN expr
```

The RETURN statement terminates execution of a stored function and returns the value *expr* to the function caller. There must be at least one RETURN statement in a stored function. There may be more than one if the function has multiple exit points.

This statement is not used in stored procedures, triggers, or events. The LEAVE statement can be used to exit a stored program of those types.

13.6.5.8 WHILE Syntax

```
[begin_label:] WHILE search_condition DO
    statement_list
END WHILE [end_label]
```

The statement list within a WHILE statement is repeated as long as the *search_condition* expression is true. *statement_list* consists of one or more SQL statements, each terminated by a semicolon (;) statement delimiter.

A WHILE statement can be labeled. For the rules regarding label use, see Section 13.6.2, “Statement Label Syntax”.

Example:

```
CREATE PROCEDURE dowhile()
BEGIN
  DECLARE v1 INT DEFAULT 5;

  WHILE v1 > 0 DO
    ...
    SET v1 = v1 - 1;
  END WHILE;
END;
```

13.6.6 Cursors

MySQL supports cursors inside stored programs. The syntax is as in embedded SQL. Cursors have these properties:

- Asensitive: The server may or may not make a copy of its result table
- Read only: Not updatable
- Nonscrollable: Can be traversed only in one direction and cannot skip rows

Cursor declarations must appear before handler declarations and after variable and condition declarations.

Example:

```
CREATE PROCEDURE curdemo()
BEGIN
  DECLARE done INT DEFAULT FALSE;
  DECLARE a CHAR(16);
  DECLARE b, c INT;
  DECLARE cur1 CURSOR FOR SELECT id,data FROM test.t1;
  DECLARE cur2 CURSOR FOR SELECT i FROM test.t2;
  DECLARE CONTINUE HANDLER FOR NOT FOUND SET done = TRUE;

  OPEN cur1;
  OPEN cur2;

  read_loop: LOOP
    FETCH cur1 INTO a, b;
    FETCH cur2 INTO c;
    IF done THEN
      LEAVE read_loop;
    END IF;
    IF b < c THEN
      INSERT INTO test.t3 VALUES (a,b);
    ELSE
      INSERT INTO test.t3 VALUES (a,c);
    END IF;
  END LOOP;

  CLOSE cur1;
  CLOSE cur2;
END;
```

13.6.6.1 Cursor CLOSE Syntax

```
CLOSE cursor_name
```

This statement closes a previously opened cursor. For an example, see Section 13.6.6, “Cursors”.

An error occurs if the cursor is not open.

If not closed explicitly, a cursor is closed at the end of the `BEGIN ... END` block in which it was declared.

13.6.6.2 Cursor DECLARE Syntax

```
DECLARE cursor_name CURSOR FOR select_statement
```

This statement declares a cursor and associates it with a `SELECT` statement that retrieves the rows to be traversed by the cursor. To fetch the rows later, use a `FETCH` statement. The number of columns retrieved by the `SELECT` statement must match the number of output variables specified in the `FETCH` statement.

The SELECT statement cannot have an INTO clause.

Cursor declarations must appear before handler declarations and after variable and condition declarations.

A stored program may contain multiple cursor declarations, but each cursor declared in a given block must have a unique name. For an example, see Section 13.6.6, “Cursors”.

For information available through SHOW statements, it is possible in many cases to obtain equivalent information by using a cursor with an INFORMATION_SCHEMA table.

13.6.6.3 Cursor FETCH Syntax

```
FETCH [[NEXT] FROM] cursor_name INTO var_name [, var_name] ...
```

This statement fetches the next row for the SELECT statement associated with the specified cursor (which must be open), and advances the cursor pointer. If a row exists, the fetched columns are stored in the named variables. The number of columns retrieved by the SELECT statement must match the number of output variables specified in the FETCH statement.

If no more rows are available, a No Data condition occurs with SQLSTATE value '02000'. To detect this condition, you can set up a handler for it (or for a NOT FOUND condition). For an example, see Section 13.6.6, “Cursors”.

Be aware that another operation, such as a SELECT or another FETCH, may also cause the handler to execute by raising the same condition. If it is necessary to distinguish which operation raised the condition, place the operation within its own BEGIN ... END block so that it can be associated with its own handler.

13.6.6.4 Cursor OPEN Syntax

```
OPEN cursor_name
```

This statement opens a previously declared cursor. For an example, see Section 13.6.6, “Cursors”.

13.6.7 Condition Handling

Conditions may arise during stored program execution that require special handling, such as exiting the current program block or continuing execution. Handlers can be defined for general conditions such as warnings or exceptions, or for specific conditions such as a particular error code. Specific conditions can be assigned names and referred to that way in handlers.

To name a condition, use the DECLARE ... CONDITION statement. To declare a handler, use the DECLARE ... HANDLER statement. See Section 13.6.7.1, “DECLARE ... CONDITION Syntax”, and Section 13.6.7.2, “DECLARE ... HANDLER Syntax”. For information about how the server chooses handlers when a condition occurs, see Section 13.6.7.6, “Scope Rules for Handlers”.

To raise a condition, use the SIGNAL statement. To modify condition information within a condition handler, use RESIGNAL. See Section 13.6.7.1, “DECLARE ... CONDITION Syntax”, and Section 13.6.7.2, “DECLARE ... HANDLER Syntax”.

To retrieve information from the diagnostics area, use the GET DIAGNOSTICS statement (see Section 13.6.7.3, “GET DIAGNOSTICS Syntax”). For information about the diagnostics area, see Section 13.6.7.7, “The MySQL Diagnostics Area”.

13.6.7.1 DECLARE ... CONDITION Syntax

```
DECLARE condition_name CONDITION FOR condition_value
```

```
condition_value:
    mysql_error_code
  | SQLSTATE [VALUE] sqlstate_value
```

The `DECLARE ... CONDITION` statement declares a named error condition, associating a name with a condition that needs specific handling. The name can be referred to in a subsequent `DECLARE ... HANDLER` statement (see Section 13.6.7.2, “DECLARE ... HANDLER Syntax”).

Condition declarations must appear before cursor or handler declarations.

The *`condition_value`* for `DECLARE ... CONDITION` indicates the specific condition or class of conditions to associate with the condition name. It can take the following forms:

- *`mysql_error_code`*: An integer literal indicating a MySQL error code.

  Do not use MySQL error code 0 because that indicates success rather than an error condition. For a list of MySQL error codes, see Section B.3, “Server Error Codes and Messages”.

- SQLSTATE [VALUE] *`sqlstate_value`*: A 5-character string literal indicating an SQLSTATE value.

  Do not use SQLSTATE values that begin with `'00'` because those indicate success rather than an error condition. For a list of SQLSTATE values, see Section B.3, “Server Error Codes and Messages”.

Condition names referred to in `SIGNAL` or use `RESIGNAL` statements must be associated with SQLSTATE values, not MySQL error codes.

Using names for conditions can help make stored program code clearer. For example, this handler applies to attempts to drop a nonexistent table, but that is apparent only if you know that 1051 is the MySQL error code for “unknown table”:

```
DECLARE CONTINUE HANDLER FOR 1051
  BEGIN
    -- body of handler
  END;
```

By declaring a name for the condition, the purpose of the handler is more readily seen:

```
DECLARE no_such_table CONDITION FOR 1051;
DECLARE CONTINUE HANDLER FOR no_such_table
  BEGIN
    -- body of handler
  END;
```

Here is a named condition for the same condition, but based on the corresponding SQLSTATE value rather than the MySQL error code:

```
DECLARE no_such_table CONDITION FOR SQLSTATE '42S02';
DECLARE CONTINUE HANDLER FOR no_such_table
  BEGIN
    -- body of handler
  END;
```

13.6.7.2 DECLARE ... HANDLER Syntax

```
DECLARE handler_action HANDLER
    FOR condition_value [, condition_value] ...
    statement

handler_action:
    CONTINUE
  | EXIT
```

```
  | UNDO

condition_value:
    mysql_error_code
  | SQLSTATE [VALUE] sqlstate_value
  | condition_name
  | SQLWARNING
  | NOT FOUND
  | SQLEXCEPTION
```

The `DECLARE ... HANDLER` statement specifies a handler that deals with one or more conditions. If one of these conditions occurs, the specified *`statement`* executes. *`statement`* can be a simple statement such as `SET var_name = value`, or a compound statement written using `BEGIN` and `END` (see Section 13.6.1, “BEGIN ... END Compound-Statement Syntax”).

Handler declarations must appear after variable or condition declarations.

The *`handler_action`* value indicates what action the handler takes after execution of the handler statement:

- `CONTINUE`: Execution of the current program continues.
- `EXIT`: Execution terminates for the `BEGIN ... END` compound statement in which the handler is declared. This is true even if the condition occurs in an inner block.
- `UNDO`: Not supported.

The *`condition_value`* for `DECLARE ... HANDLER` indicates the specific condition or class of conditions that activates the handler. It can take the following forms:

- *`mysql_error_code`*: An integer literal indicating a MySQL error code, such as 1051 to specify “unknown table”:

  ```
  DECLARE CONTINUE HANDLER FOR 1051
    BEGIN
      -- body of handler
    END;
  ```

  Do not use MySQL error code 0 because that indicates success rather than an error condition. For a list of MySQL error codes, see Section B.3, “Server Error Codes and Messages”.

- SQLSTATE [VALUE] *`sqlstate_value`*: A 5-character string literal indicating an SQLSTATE value, such as `'42S01'` to specify “unknown table”:

  ```
  DECLARE CONTINUE HANDLER FOR SQLSTATE '42S02'
    BEGIN
      -- body of handler
    END;
  ```

  Do not use SQLSTATE values that begin with `'00'` because those indicate success rather than an error condition. For a list of SQLSTATE values, see Section B.3, “Server Error Codes and Messages”.

- *`condition_name`*: A condition name previously specified with `DECLARE ... CONDITION`. A condition name can be associated with a MySQL error code or SQLSTATE value. See Section 13.6.7.1, “DECLARE ... CONDITION Syntax”.

- `SQLWARNING`: Shorthand for the class of SQLSTATE values that begin with `'01'`.

  ```
  DECLARE CONTINUE HANDLER FOR SQLWARNING
    BEGIN
      -- body of handler
    END;
  ```

- `NOT FOUND`: Shorthand for the class of SQLSTATE values that begin with `'02'`. This is relevant within the context of cursors and is used to control what happens when a cursor reaches the end of a data set. If no more rows are available, a No Data condition occurs with SQLSTATE value `'02000'`. To detect this condition, you can set up a handler for it or for a `NOT FOUND` condition.

```
DECLARE CONTINUE HANDLER FOR NOT FOUND
  BEGIN
    -- body of handler
  END;
```

  For another example, see Section 13.6.6, “Cursors”. The `NOT FOUND` condition also occurs for `SELECT ... INTO var_list` statements that retrieve no rows.

- `SQLEXCEPTION`: Shorthand for the class of SQLSTATE values that do not begin with `'00'`, `'01'`, or `'02'`.

```
DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
  BEGIN
    -- body of handler
  END;
```

For information about how the server chooses handlers when a condition occurs, see Section 13.6.7.6, “Scope Rules for Handlers”.

If a condition occurs for which no handler has been declared, the action taken depends on the condition class:

- For `SQLEXCEPTION` conditions, the stored program terminates at the statement that raised the condition, as if there were an `EXIT` handler. If the program was called by another stored program, the calling program handles the condition using the handler selection rules applied to its own handlers.

- For `SQLWARNING` conditions, the program continues executing, as if there were a `CONTINUE` handler.

- For `NOT FOUND` conditions, if the condition was raised normally, the action is `CONTINUE`. If it was raised by `SIGNAL` or `RESIGNAL`, the action is `EXIT`.

The following example uses a handler for `SQLSTATE '23000'`, which occurs for a duplicate-key error:

```
mysql> CREATE TABLE test.t (s1 INT, PRIMARY KEY (s1));
Query OK, 0 rows affected (0.00 sec)

mysql> delimiter //

mysql> CREATE PROCEDURE handlerdemo ()
    -> BEGIN
    ->   DECLARE CONTINUE HANDLER FOR SQLSTATE '23000' SET @x2 = 1;
    ->   SET @x = 1;
    ->   INSERT INTO test.t VALUES (1);
    ->   SET @x = 2;
    ->   INSERT INTO test.t VALUES (1);
    ->   SET @x = 3;
    -> END;
    -> //
Query OK, 0 rows affected (0.00 sec)

mysql> CALL handlerdemo()//
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @x//
    +------+
    | @x   |
    +------+
```

```
    | 3    |
    +------+
    1 row in set (0.00 sec)
```

Notice that `@x` is `3` after the procedure executes, which shows that execution continued to the end of the procedure after the error occurred. If the `DECLARE ... HANDLER` statement had not been present, MySQL would have taken the default action (`EXIT`) after the second `INSERT` failed due to the `PRIMARY KEY` constraint, and `SELECT @x` would have returned `2`.

To ignore a condition, declare a `CONTINUE` handler for it and associate it with an empty block. For example:

```
DECLARE CONTINUE HANDLER FOR SQLWARNING BEGIN END;
```

The scope of a block label does not include the code for handlers declared within the block. Therefore, the statement associated with a handler cannot use `ITERATE` or `LEAVE` to refer to labels for blocks that enclose the handler declaration. Consider the following example, where the `REPEAT` block has a label of `retry`:

```
CREATE PROCEDURE p ()
BEGIN
  DECLARE i INT DEFAULT 3;
  retry:
    REPEAT
      BEGIN
        DECLARE CONTINUE HANDLER FOR SQLWARNING
          BEGIN
            ITERATE retry;    # illegal
          END;
        IF i < 0 THEN
          LEAVE retry;        # legal
        END IF;
        SET i = i - 1;
      END;
    UNTIL FALSE END REPEAT;
END;
```

The `retry` label is in scope for the `IF` statement within the block. It is not in scope for the `CONTINUE` handler, so the reference there is invalid and results in an error:

```
ERROR 1308 (42000): LEAVE with no matching label: retry
```

To avoid references to outer labels in handlers, use one of these strategies:

- To leave the block, use an `EXIT` handler. If no block cleanup is required, the `BEGIN ... END` handler body can be empty:

  ```
  DECLARE EXIT HANDLER FOR SQLWARNING BEGIN END;
  ```

  Otherwise, put the cleanup statements in the handler body:

  ```
  DECLARE EXIT HANDLER FOR SQLWARNING
    BEGIN
      block cleanup statements
    END;
  ```

- To continue execution, set a status variable in a `CONTINUE` handler that can be checked in the enclosing block to determine whether the handler was invoked. The following example uses the variable `done` for this purpose:

  ```
  CREATE PROCEDURE p ()
  BEGIN
  ```

```
  DECLARE i INT DEFAULT 3;
  DECLARE done INT DEFAULT FALSE;
  retry:
    REPEAT
      BEGIN
        DECLARE CONTINUE HANDLER FOR SQLWARNING
          BEGIN
            SET done = TRUE;
          END;
        IF done OR i < 0 THEN
          LEAVE retry;
        END IF;
        SET i = i - 1;
      END;
    UNTIL FALSE END REPEAT;
END;
```

13.6.7.3 GET DIAGNOSTICS Syntax

```
GET [CURRENT | STACKED] DIAGNOSTICS
{
    statement_information_item
    [, statement_information_item] ...
  | CONDITION condition_number
    condition_information_item
    [, condition_information_item] ...
}

statement_information_item:
    target = statement_information_item_name

condition_information_item:
    target = condition_information_item_name

statement_information_item_name:
    NUMBER
  | ROW_COUNT

condition_information_item_name:
    CLASS_ORIGIN
  | SUBCLASS_ORIGIN
  | RETURNED_SQLSTATE
  | MESSAGE_TEXT
  | MYSQL_ERRNO
  | CONSTRAINT_CATALOG
  | CONSTRAINT_SCHEMA
  | CONSTRAINT_NAME
  | CATALOG_NAME
  | SCHEMA_NAME
  | TABLE_NAME
  | COLUMN_NAME
  | CURSOR_NAME

condition_number, target:
    (see following discussion)
```

SQL statements produce diagnostic information that populates the diagnostics area. The `GET DIAGNOSTICS` statement enables applications to inspect this information. (You can also use `SHOW WARNINGS` or `SHOW ERRORS` to see conditions or errors.)

No special privileges are required to execute `GET DIAGNOSTICS`.

The keyword `CURRENT` means to retrieve information from the current diagnostics area. The keyword `STACKED` means to retrieve information from the second diagnostics area, which is available only if the current context is a condition handler. If neither keyword is given, the default is to use the current diagnostics area.

The `GET DIAGNOSTICS` statement is typically used in a handler within a stored program. It is a MySQL extension that `GET [CURRENT] DIAGNOSTICS` is permitted outside handler context to check

the execution of any SQL statement. For example, if you invoke the `mysql` client program, you can enter these statements at the prompt:

```
mysql> DROP TABLE test.no_such_table;
ERROR 1051 (42S02): Unknown table 'test.no_such_table'
mysql> GET DIAGNOSTICS CONDITION 1
    ->   @p1 = RETURNED_SQLSTATE, @p2 = MESSAGE_TEXT;
mysql> SELECT @p1, @p2;
+-------+-----------------------------------+
| @p1   | @p2                               |
+-------+-----------------------------------+
| 42S02 | Unknown table 'test.no_such_table' |
+-------+-----------------------------------+
```

This extension applies only to the current diagnostics area. It does not apply to the second diagnostics area because `GET STACKED DIAGNOSTICS` is permitted only if the current context is a condition handler. If that is not the case, a `GET STACKED DIAGNOSTICS when handler not active` error occurs.

For a description of the diagnostics area, see Section 13.6.7.7, “The MySQL Diagnostics Area”. Briefly, it contains two kinds of information:

- Statement information, such as the number of conditions that occurred or the affected-rows count.
- Condition information, such as the error code and message. If a statement raises multiple conditions, this part of the diagnostics area has a condition area for each one. If a statement raises no conditions, this part of the diagnostics area is empty.

For a statement that produces three conditions, the diagnostics area contains statement and condition information like this:

```
Statement information:
  row count
  ... other statement information items ...
Condition area list:
  Condition area 1:
    error code for condition 1
    error message for condition 1
    ... other condition information items ...
  Condition area 2:
    error code for condition 2:
    error message for condition 2
    ... other condition information items ...
  Condition area 3:
    error code for condition 3
    error message for condition 3
    ... other condition information items ...
```

`GET DIAGNOSTICS` can obtain either statement or condition information, but not both in the same statement:

- To obtain statement information, retrieve the desired statement items into target variables. This instance of `GET DIAGNOSTICS` assigns the number of available conditions and the rows-affected count to the user variables `@p1` and `@p2`:

  ```
  GET DIAGNOSTICS @p1 = NUMBER, @p2 = ROW_COUNT;
  ```

- To obtain condition information, specify the condition number and retrieve the desired condition items into target variables. This instance of `GET DIAGNOSTICS` assigns the SQLSTATE value and error message to the user variables `@p3` and `@p4`:

  ```
  GET DIAGNOSTICS CONDITION 1
  ```

```
  @p3 = RETURNED_SQLSTATE, @p4 = MESSAGE_TEXT;
```

The retrieval list specifies one or more *`target = item_name`* assignments, separated by commas. Each assignment names a target variable and either a *`statement_information_item_name`* or *`condition_information_item_name`* designator, depending on whether the statement retrieves statement or condition information.

Valid *`target`* designators for storing item information can be stored procedure or function parameters, stored program local variables declared with `DECLARE`, or user-defined variables.

Valid *`condition_number`* designators can be stored procedure or function parameters, stored program local variables declared with `DECLARE`, user-defined variables, system variables, or literals. A character literal may include a *`_charset`* introducer. A warning occurs if the condition number is not in the range from 1 to the number of condition areas that have information. In this case, the warning is added to the diagnostics area without clearing it.

When a condition occurs, MySQL does not populate all condition items recognized by `GET DIAGNOSTICS`. For example:

```
mysql> GET DIAGNOSTICS CONDITION 1
    ->   @p5 = SCHEMA_NAME, @p6 = TABLE_NAME;
mysql> SELECT @p5, @p6;
+------+------+
| @p5  | @p6  |
+------+------+
|      |      |
+------+------+
```

In standard SQL, if there are multiple conditions, the first condition relates to the `SQLSTATE` value returned for the previous SQL statement. In MySQL, this is not guaranteed. To get the main error, you cannot do this:

```
GET DIAGNOSTICS CONDITION 1 @errno = MYSQL_ERRNO;
```

Instead, retrieve the condition count first, then use it to specify which condition number to inspect:

```
GET DIAGNOSTICS @cno = NUMBER;
GET DIAGNOSTICS CONDITION @cno @errno = MYSQL_ERRNO;
```

For information about permissible statement and condition information items, and which ones are populated when a condition occurs, see Diagnostics Area Information Items.

Here is an example that uses `GET DIAGNOSTICS` and an exception handler in stored procedure context to assess the outcome of an insert operation. If the insert was successful, the procedure uses `GET DIAGNOSTICS` to get the rows-affected count. This shows that you can use `GET DIAGNOSTICS` multiple times to retrieve information about a statement as long as the current diagnostics area has not been cleared.

```
CREATE PROCEDURE do_insert(value INT)
BEGIN
  -- Declare variables to hold diagnostics area information
  DECLARE code CHAR(5) DEFAULT '00000';
  DECLARE msg TEXT;
  DECLARE rows INT;
  DECLARE result TEXT;
  -- Declare exception handler for failed insert
  DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
    BEGIN
      GET DIAGNOSTICS CONDITION 1
        code = RETURNED_SQLSTATE, msg = MESSAGE_TEXT;
    END;
```

```
  -- Perform the insert
  INSERT INTO t1 (int_col) VALUES(value);
  -- Check whether the insert was successful
  IF code = '00000' THEN
    GET DIAGNOSTICS rows = ROW_COUNT;
    SET result = CONCAT('insert succeeded, row count = ',rows);
  ELSE
    SET result = CONCAT('insert failed, error = ',code,', message = ',msg);
  END IF;
  -- Say what happened
  SELECT result;
END;
```

Suppose that `t1.int_col` is an integer column that is declared as `NOT NULL`. The procedure produces these results when invoked to insert non-`NULL` and `NULL` values, respectively:

```
mysql> CALL do_insert(1);
+---------------------------------+
| result                          |
+---------------------------------+
| insert succeeded, row count = 1 |
+---------------------------------+

mysql> CALL do_insert(NULL);
+-------------------------------------------------------------------------+
| result                                                                  |
+-------------------------------------------------------------------------+
| insert failed, error = 23000, message = Column 'int_col' cannot be null |
+-------------------------------------------------------------------------+
```

When a condition handler activates, a push to the diagnostics area stack occurs:

- The first (current) diagnostics area becomes the second (stacked) diagnostics area and a new current diagnostics area is created as a copy of it.
- `GET [CURRENT] DIAGNOSTICS` and `GET STACKED DIAGNOSTICS` can be used within the handler to access the contents of the current and stacked diagnostics areas.
- Initially, both diagnostics areas return the same result, so it is possible to get information from the current diagnostics area about the condition that activated the handler, *as long as* you execute no statements within the handler that change its current diagnostics area.
- However, statements executing within the handler can modify the current diagnostics area, clearing and setting its contents according to the normal rules (see How the Diagnostics Area is Populated).

  A more reliable way to obtain information about the handler-activating condition is to use the stacked diagnostics area, which cannot be modified by statements executing within the handler except `RESIGNAL`. For information about when the current diagnostics area is set and cleared, see Section 13.6.7.7, “The MySQL Diagnostics Area”.

The next example shows how `GET STACKED DIAGNOSTICS` can be used within a handler to obtain information about the handled exception, even after the current diagnostics area has been modified by handler statements.

Within a stored procedure `p()`, we attempt to insert two values into a table that contains a `TEXT NOT NULL` column. The first value is a non-`NULL` string and the second is `NULL`. The column prohibits `NULL` values, so the first insert succeeds but the second causes an exception. The procedure includes an exception handler that maps attempts to insert `NULL` into inserts of the empty string:

```
DROP TABLE IF EXISTS t1;
CREATE TABLE t1 (c1 TEXT NOT NULL);
DROP PROCEDURE IF EXISTS p;
delimiter //
```

```
CREATE PROCEDURE p ()
BEGIN
  -- Declare variables to hold diagnostics area information
  DECLARE errcount INT;
  DECLARE errno INT;
  DECLARE msg TEXT;
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    -- Here the current DA is nonempty because no prior statements
    -- executing within the handler have cleared it
    GET CURRENT DIAGNOSTICS CONDITION 1
      errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
    SELECT 'current DA before mapped insert' AS op, errno, msg;
    GET STACKED DIAGNOSTICS CONDITION 1
      errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
    SELECT 'stacked DA before mapped insert' AS op, errno, msg;

    -- Map attempted NULL insert to empty string insert
    INSERT INTO t1 (c1) VALUES('');

    -- Here the current DA should be empty (if the INSERT succeeded),
    -- so check whether there are conditions before attempting to
    -- obtain condition information
    GET CURRENT DIAGNOSTICS errcount = NUMBER;
    IF errcount = 0
    THEN
      SELECT 'mapped insert succeeded, current DA is empty' AS op;
    ELSE
      GET CURRENT DIAGNOSTICS CONDITION 1
        errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
      SELECT 'current DA after mapped insert' AS op, errno, msg;
    END IF ;
    GET STACKED DIAGNOSTICS CONDITION 1
      errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
    SELECT 'stacked DA after mapped insert' AS op, errno, msg;
  END;
  INSERT INTO t1 (c1) VALUES('string 1');
  INSERT INTO t1 (c1) VALUES(NULL);
END;
//
delimiter ;
CALL p();
SELECT * FROM t1;
```

When the handler activates, a copy of the current diagnostics area is pushed to the diagnostics area stack. The handler first displays the contents of the current and stacked diagnostics areas, which are both the same initially:

```
+---------------------------------+-------+----------------------------+
| op                              | errno | msg                        |
+---------------------------------+-------+----------------------------+
| current DA before mapped insert |  1048 | Column 'c1' cannot be null |
+---------------------------------+-------+----------------------------+

+---------------------------------+-------+----------------------------+
| op                              | errno | msg                        |
+---------------------------------+-------+----------------------------+
| stacked DA before mapped insert |  1048 | Column 'c1' cannot be null |
+---------------------------------+-------+----------------------------+
```

Statements executing after the `GET DIAGNOSTICS` statements may reset the current diagnostics area. statements may reset the current diagnostics area. For example, the handler maps the `NULL` insert to an empty-string insert and displays the result. The new insert succeeds and clears the current diagnostics area, but the stacked diagnostics area remains unchanged and still contains information about the condition that activated the handler:

```
+--------------------------------------------+
| op                                         |
```

```
+-------------------------------------------+
| mapped insert succeeded, current DA is empty |
+-------------------------------------------+

+--------------------------------+-------+----------------------------+
| op                             | errno | msg                        |
+--------------------------------+-------+----------------------------+
| stacked DA after mapped insert |  1048 | Column 'c1' cannot be null |
+--------------------------------+-------+----------------------------+
```

When the condition handler ends, its current diagnostics area is popped from the stack and the stacked diagnostics area becomes the current diagnostics area in the stored procedure.

After the procedure returns, the table contains two rows. The empty row results from the attempt to insert `NULL` that was mapped to an empty-string insert:

```
+----------+
| c1       |
+----------+
| string 1 |
|          |
+----------+
```

In the preceding example, the first two `GET DIAGNOSTICS` statements within the condition handler that retrieve information from the current and stacked diagnostics areas return the same values. This will not be the case if statements that reset the current diagnostics area execute earlier within the handler. Suppose that `p()` is rewritten to place the `DECLARE` statements within the handler definition rather than preceding it:

```
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    -- Declare variables to hold diagnostics area information
    DECLARE errcount INT;
    DECLARE errno INT;
    DECLARE msg TEXT;
    GET CURRENT DIAGNOSTICS CONDITION 1
      errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
    SELECT 'current DA before mapped insert' AS op, errno, msg;
    GET STACKED DIAGNOSTICS CONDITION 1
      errno = MYSQL_ERRNO, msg = MESSAGE_TEXT;
    SELECT 'stacked DA before mapped insert' AS op, errno, msg;
...
```

In this case, the result is version dependent:

- Before MySQL 5.7.2, `DECLARE` does not change the current diagnostics area, so the first two `GET DIAGNOSTICS` statements return the same result, just as in the original version of `p()`.

  In MySQL 5.7.2, work was done to ensure that all nondiagnostic statements populate the diagnostics area, per the SQL standard. `DECLARE` is one of them, so in 5.7.2 and higher, `DECLARE` statements executing at the beginning of the handler clear the current diagnostics area and the `GET DIAGNOSTICS` statements produce different results:

  ```
  +--------------------------------+-------+------+
  | op                             | errno | msg  |
  +--------------------------------+-------+------+
  | current DA before mapped insert |  NULL | NULL |
  +--------------------------------+-------+------+

  +--------------------------------+-------+----------------------------+
  | op                             | errno | msg                        |
  +--------------------------------+-------+----------------------------+
  | stacked DA before mapped insert |  1048 | Column 'c1' cannot be null |
  ```

```
+--------------------------------+-------+--------------------------+
```

To avoid this issue within a condition handler when seeking to obtain information about the condition that activated the handler, be sure to access the stacked diagnostics area, not the current diagnostics area.

13.6.7.4 RESIGNAL Syntax

```
RESIGNAL [condition_value]
    [SET signal_information_item
    [, signal_information_item] ...]

condition_value:
    SQLSTATE [VALUE] sqlstate_value
  | condition_name

signal_information_item:
    condition_information_item_name = simple_value_specification

condition_information_item_name:
    CLASS_ORIGIN
  | SUBCLASS_ORIGIN
  | MESSAGE_TEXT
  | MYSQL_ERRNO
  | CONSTRAINT_CATALOG
  | CONSTRAINT_SCHEMA
  | CONSTRAINT_NAME
  | CATALOG_NAME
  | SCHEMA_NAME
  | TABLE_NAME
  | COLUMN_NAME
  | CURSOR_NAME

condition_name, simple_value_specification:
    (see following discussion)
```

`RESIGNAL` passes on the error condition information that is available during execution of a condition handler within a compound statement inside a stored procedure or function, trigger, or event. `RESIGNAL` may change some or all information before passing it on. `RESIGNAL` is related to `SIGNAL`, but instead of originating a condition as `SIGNAL` does, `RESIGNAL` relays existing condition information, possibly after modifying it.

`RESIGNAL` makes it possible to both handle an error and return the error information. Otherwise, by executing an SQL statement within the handler, information that caused the handler's activation is destroyed. `RESIGNAL` also can make some procedures shorter if a given handler can handle part of a situation, then pass the condition “up the line” to another handler.

No special privileges are required to execute the `RESIGNAL` statement.

All forms of `RESIGNAL` require that the current context be a condition handler. Otherwise, `RESIGNAL` is illegal and a `RESIGNAL when handler not active` error occurs.

To retrieve information from the diagnostics area, use the `GET DIAGNOSTICS` statement (see Section 13.6.7.3, “GET DIAGNOSTICS Syntax”). For information about the diagnostics area, see Section 13.6.7.7, “The MySQL Diagnostics Area”.

For *`condition_value`* and *`signal_information_item`*, the definitions and rules are the same for `RESIGNAL` as for `SIGNAL`. For example, the *`condition_value`* can be an `SQLSTATE` value, and the value can indicate errors, warnings, or “not found.” For additional information, see Section 13.6.7.5, “SIGNAL Syntax”.

The `RESIGNAL` statement takes *`condition_value`* and `SET` clauses, both of which are optional. This leads to several possible uses:

- `RESIGNAL` alone:

```
RESIGNAL;
```

- `RESIGNAL` with new signal information:

```
RESIGNAL SET signal_information_item [, signal_information_item] ...;
```

- `RESIGNAL` with a condition value and possibly new signal information:

```
RESIGNAL condition_value
    [SET signal_information_item [, signal_information_item] ...];
```

These use cases all cause changes to the diagnostics and condition areas:

- A diagnostics area contains one or more condition areas.
- A condition area contains condition information items, such as the `SQLSTATE` value, `MYSQL_ERRNO`, or `MESSAGE_TEXT`.

There is a stack of diagnostics areas. When a handler takes control, it pushes a diagnostics area to the top of the stack, so there are two diagnostics areas during handler execution:

- The first (current) diagnostics area, which starts as a copy of the last diagnostics area, but will be overwritten by the first statement in the handler that changes the current diagnostics area.
- The last (stacked) diagnostics area, which has the condition areas that were set up before the handler took control.

The maximum number of condition areas in a diagnostics area is determined by the value of the `max_error_count` system variable. See Diagnostics Area-Related System Variables.

RESIGNAL Alone

A simple `RESIGNAL` alone means “pass on the error with no change.” It restores the last diagnostics area and makes it the current diagnostics area. That is, it “pops” the diagnostics area stack.

Within a condition handler that catches a condition, one use for `RESIGNAL` alone is to perform some other actions, and then pass on without change the original condition information (the information that existed before entry into the handler).

Example:

```
DROP TABLE IF EXISTS xx;
delimiter //
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    SET @error_count = @error_count + 1;
    IF @a = 0 THEN RESIGNAL; END IF;
  END;
  DROP TABLE xx;
END//
delimiter ;
SET @error_count = 0;
SET @a = 0;
CALL p();
```

Suppose that the `DROP TABLE xx` statement fails. The diagnostics area stack looks like this:

```
DA 1. ERROR 1051 (42S02): Unknown table 'xx'
```

Then execution enters the `EXIT` handler. It starts by pushing a diagnostics area to the top of the stack, which now looks like this:

```
DA 1. ERROR 1051 (42S02): Unknown table 'xx'
DA 2. ERROR 1051 (42S02): Unknown table 'xx'
```

At this point, the contents of the first (current) and second (stacked) diagnostics areas are the same. The first diagnostics area may be modified by statements executing subsequently within the handler.

Usually a procedure statement clears the first diagnostics area. `BEGIN` is an exception, it does not clear, it does nothing. `SET` is not an exception, it clears, performs the operation, and produces a result of “success.” The diagnostics area stack now looks like this:

```
DA 1. ERROR 0000 (00000): Successful operation
DA 2. ERROR 1051 (42S02): Unknown table 'xx'
```

At this point, if `@a = 0`, `RESIGNAL` pops the diagnostics area stack, which now looks like this:

```
DA 1. ERROR 1051 (42S02): Unknown table 'xx'
```

And that is what the caller sees.

If `@a` is not 0, the handler simply ends, which means that there is no more use for the current diagnostics area (it has been “handled”), so it can be thrown away, causing the stacked diagnostics area to become the current diagnostics area again. The diagnostics area stack looks like this:

```
DA 1. ERROR 0000 (00000): Successful operation
```

The details make it look complex, but the end result is quite useful: Handlers can execute without destroying information about the condition that caused activation of the handler.

RESIGNAL with New Signal Information

`RESIGNAL` with a `SET` clause provides new signal information, so the statement means “pass on the error with changes”:

```
RESIGNAL SET signal_information_item [, signal_information_item] ...;
```

As with `RESIGNAL` alone, the idea is to pop the diagnostics area stack so that the original information will go out. Unlike `RESIGNAL` alone, anything specified in the `SET` clause changes.

Example:

```
DROP TABLE IF EXISTS xx;
delimiter //
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    SET @error_count = @error_count + 1;
    IF @a = 0 THEN RESIGNAL SET MYSQL_ERRNO = 5; END IF;
  END;
  DROP TABLE xx;
END//
delimiter ;
SET @error_count = 0;
SET @a = 0;
```

```
CALL p();
```

Remember from the previous discussion that `RESIGNAL` alone results in a diagnostics area stack like this:

```
DA 1. ERROR 1051 (42S02): Unknown table 'xx'
```

The `RESIGNAL SET MYSQL_ERRNO = 5` statement results in this stack instead, which is what the caller sees:

```
DA 1. ERROR 5 (42S02): Unknown table 'xx'
```

In other words, it changes the error number, and nothing else.

The `RESIGNAL` statement can change any or all of the signal information items, making the first condition area of the diagnostics area look quite different.

RESIGNAL with a Condition Value and Optional New Signal Information

`RESIGNAL` with a condition value means “push a condition into the current diagnostics area.” If the `SET` clause is present, it also changes the error information.

```
RESIGNAL condition_value
    [SET signal_information_item [, signal_information_item] ...];
```

This form of `RESIGNAL` restores the last diagnostics area and makes it the current diagnostics area. That is, it “pops” the diagnostics area stack, which is the same as what a simple `RESIGNAL` alone would do. However, it also changes the diagnostics area depending on the condition value or signal information.

Example:

```
DROP TABLE IF EXISTS xx;
delimiter //
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    SET @error_count = @error_count + 1;
    IF @a = 0 THEN RESIGNAL SQLSTATE '45000' SET MYSQL_ERRNO=5; END IF;
  END;
  DROP TABLE xx;
END//
delimiter ;
SET @error_count = 0;
SET @a = 0;
SET @@max_error_count = 2;
CALL p();
SHOW ERRORS;
```

This is similar to the previous example, and the effects are the same, except that if `RESIGNAL` happens, the current condition area looks different at the end. (The reason the condition adds to rather than replaces the existing condition is the use of a condition value.)

The `RESIGNAL` statement includes a condition value (`SQLSTATE '45000'`), so it adds a new condition area, resulting in a diagnostics area stack that looks like this:

```
DA 1. (condition 2) ERROR 1051 (42S02): Unknown table 'xx'
      (condition 1) ERROR 5 (45000) Unknown table 'xx'
```

The result of `CALL p()` and `SHOW ERRORS` for this example is:

```
mysql> CALL p();
ERROR 5 (45000): Unknown table 'xx'
mysql> SHOW ERRORS;
+-------+------+---------------------------------+
| Level | Code | Message                         |
+-------+------+---------------------------------+
| Error | 1051 | Unknown table 'xx'              |
| Error |    5 | Unknown table 'xx'              |
+-------+------+---------------------------------+
```

RESIGNAL Requires Condition Handler Context

All forms of `RESIGNAL` require that the current context be a condition handler. Otherwise, `RESIGNAL` is illegal and a `RESIGNAL when handler not active` error occurs. For example:

```
mysql> CREATE PROCEDURE p () RESIGNAL;
Query OK, 0 rows affected (0.00 sec)

mysql> CALL p();
ERROR 1645 (0K000): RESIGNAL when handler not active
```

Here is a more difficult example:

```
delimiter //
CREATE FUNCTION f () RETURNS INT
BEGIN
  RESIGNAL;
  RETURN 5;
END//
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION SET @a=f();
  SIGNAL SQLSTATE '55555';
END//
delimiter ;
CALL p();
```

`RESIGNAL` occurs within the stored function `f()`. Although `f()` itself is invoked within the context of the `EXIT` handler, execution within `f()` has its own context, which is not handler context. Thus, `RESIGNAL` within `f()` results in a “handler not active” error.

13.6.7.5 SIGNAL Syntax

```
SIGNAL condition_value
    [SET signal_information_item
    [, signal_information_item] ...]

condition_value:
    SQLSTATE [VALUE] sqlstate_value
  | condition_name

signal_information_item:
    condition_information_item_name = simple_value_specification

condition_information_item_name:
    CLASS_ORIGIN
  | SUBCLASS_ORIGIN
  | MESSAGE_TEXT
  | MYSQL_ERRNO
  | CONSTRAINT_CATALOG
  | CONSTRAINT_SCHEMA
  | CONSTRAINT_NAME
  | CATALOG_NAME
  | SCHEMA_NAME
  | TABLE_NAME
```

```
  | COLUMN_NAME
  | CURSOR_NAME

condition_name, simple_value_specification:
    (see following discussion)
```

`SIGNAL` is the way to “return” an error. `SIGNAL` provides error information to a handler, to an outer portion of the application, or to the client. Also, it provides control over the error's characteristics (error number, `SQLSTATE` value, message). Without `SIGNAL`, it is necessary to resort to workarounds such as deliberately referring to a nonexistent table to cause a routine to return an error.

No special privileges are required to execute the `SIGNAL` statement.

To retrieve information from the diagnostics area, use the `GET DIAGNOSTICS` statement (see Section 13.6.7.3, “GET DIAGNOSTICS Syntax”). For information about the diagnostics area, see Section 13.6.7.7, “The MySQL Diagnostics Area”.

The *`condition_value`* in a `SIGNAL` statement indicates the error value to be returned. It can be an `SQLSTATE` value (a 5-character string literal) or a *`condition_name`* that refers to a named condition previously defined with `DECLARE ... CONDITION` (see Section 13.6.7.1, “DECLARE ... CONDITION Syntax”).

An `SQLSTATE` value can indicate errors, warnings, or “not found.” The first two characters of the value indicate its error class, as discussed in Signal Condition Information Items. Some signal values cause statement termination; see Effect of Signals on Handlers, Cursors, and Statements.

The `SQLSTATE` value for a `SIGNAL` statement should not start with `'00'` because such values indicate success and are not valid for signaling an error. This is true whether the `SQLSTATE` value is specified directly in the `SIGNAL` statement or in a named condition referred to in the statement. If the value is invalid, a `Bad SQLSTATE` error occurs.

To signal a generic `SQLSTATE` value, use `'45000'`, which means “unhandled user-defined exception.”

The `SIGNAL` statement optionally includes a `SET` clause that contains multiple signal items, in a comma-separated list of *`condition_information_item_name`* = *`simple_value_specification`* assignments.

Each *`condition_information_item_name`* may be specified only once in the `SET` clause. Otherwise, a `Duplicate condition information item` error occurs.

Valid *`simple_value_specification`* designators can be specified using stored procedure or function parameters, stored program local variables declared with `DECLARE`, user-defined variables, system variables, or literals. A character literal may include a *`_charset`* introducer.

For information about permissible *`condition_information_item_name`* values, see Signal Condition Information Items.

The following procedure signals an error or warning depending on the value of `pval`, its input parameter:

```
CREATE PROCEDURE p (pval INT)
BEGIN
  DECLARE specialty CONDITION FOR SQLSTATE '45000';
  IF pval = 0 THEN
    SIGNAL SQLSTATE '01000';
  ELSEIF pval = 1 THEN
    SIGNAL SQLSTATE '45000'
      SET MESSAGE_TEXT = 'An error occurred';
  ELSEIF pval = 2 THEN
    SIGNAL specialty
      SET MESSAGE_TEXT = 'An error occurred';
  ELSE
```

```
    SIGNAL SQLSTATE '01000'
      SET MESSAGE_TEXT = 'A warning occurred', MYSQL_ERRNO = 1000;
    SIGNAL SQLSTATE '45000'
      SET MESSAGE_TEXT = 'An error occurred', MYSQL_ERRNO = 1001;
  END IF;
END;
```

If `pval` is 0, `p()` signals a warning because `SQLSTATE` values that begin with `'01'` are signals in the warning class. The warning does not terminate the procedure, and can be seen with `SHOW WARNINGS` after the procedure returns.

If `pval` is 1, `p()` signals an error and sets the `MESSAGE_TEXT` condition information item. The error terminates the procedure, and the text is returned with the error information.

If `pval` is 2, the same error is signaled, although the `SQLSTATE` value is specified using a named condition in this case.

If `pval` is anything else, `p()` first signals a warning and sets the message text and error number condition information items. This warning does not terminate the procedure, so execution continues and `p()` then signals an error. The error does terminate the procedure. The message text and error number set by the warning are replaced by the values set by the error, which are returned with the error information.

`SIGNAL` is typically used within stored programs, but it is a MySQL extension that it is permitted outside handler context. For example, if you invoke the `mysql` client program, you can enter any of these statements at the prompt:

```
mysql> SIGNAL SQLSTATE '77777';
mysql> CREATE TRIGGER t_bi BEFORE INSERT ON t
    -> FOR EACH ROW SIGNAL SQLSTATE '77777';
mysql> CREATE EVENT e ON SCHEDULE EVERY 1 SECOND
    -> DO SIGNAL SQLSTATE '77777';
```

`SIGNAL` executes according to the following rules:

If the `SIGNAL` statement indicates a particular `SQLSTATE` value, that value is used to signal the condition specified. Example:

```
CREATE PROCEDURE p (divisor INT)
BEGIN
  IF divisor = 0 THEN
    SIGNAL SQLSTATE '22012';
  END IF;
END;
```

If the `SIGNAL` statement uses a named condition, the condition must be declared in some scope that applies to the `SIGNAL` statement, and must be defined using an `SQLSTATE` value, not a MySQL error number. Example:

```
CREATE PROCEDURE p (divisor INT)
BEGIN
  DECLARE divide_by_zero CONDITION FOR SQLSTATE '22012';
  IF divisor = 0 THEN
    SIGNAL divide_by_zero;
  END IF;
END;
```

If the named condition does not exist in the scope of the `SIGNAL` statement, an `Undefined CONDITION` error occurs.

If `SIGNAL` refers to a named condition that is defined with a MySQL error number rather than an `SQLSTATE` value, a `SIGNAL/RESIGNAL can only use a CONDITION defined with`

SQLSTATE error occurs. The following statements cause that error because the named condition is associated with a MySQL error number:

```
DECLARE no_such_table CONDITION FOR 1051;
SIGNAL no_such_table;
```

If a condition with a given name is declared multiple times in different scopes, the declaration with the most local scope applies. Consider the following procedure:

```
CREATE PROCEDURE p (divisor INT)
BEGIN
  DECLARE my_error CONDITION FOR SQLSTATE '45000';
  IF divisor = 0 THEN
    BEGIN
      DECLARE my_error CONDITION FOR SQLSTATE '22012';
      SIGNAL my_error;
    END;
  END IF;
  SIGNAL my_error;
END;
```

If divisor is 0, the first SIGNAL statement executes. The innermost my_error condition declaration applies, raising SQLSTATE '22012'.

If divisor is not 0, the second SIGNAL statement executes. The outermost my_error condition declaration applies, raising SQLSTATE '45000'.

For information about how the server chooses handlers when a condition occurs, see Section 13.6.7.6, “Scope Rules for Handlers”.

Signals can be raised within exception handlers:

```
CREATE PROCEDURE p ()
BEGIN
  DECLARE EXIT HANDLER FOR SQLEXCEPTION
  BEGIN
    SIGNAL SQLSTATE VALUE '99999'
      SET MESSAGE_TEXT = 'An error occurred';
  END;
  DROP TABLE no_such_table;
END;
```

CALL p() reaches the DROP TABLE statement. There is no table named no_such_table, so the error handler is activated. The error handler destroys the original error (“no such table”) and makes a new error with SQLSTATE '99999' and message An error occurred.

Signal Condition Information Items

The following table lists the names of diagnostics area condition information items that can be set in a SIGNAL (or RESIGNAL) statement. All items are standard SQL except MYSQL_ERRNO, which is a MySQL extension. For more information about these items see Section 13.6.7.7, “The MySQL Diagnostics Area”.

```
Item Name             Definition
---------             ----------
CLASS_ORIGIN          VARCHAR(64)
SUBCLASS_ORIGIN       VARCHAR(64)
CONSTRAINT_CATALOG    VARCHAR(64)
CONSTRAINT_SCHEMA     VARCHAR(64)
CONSTRAINT_NAME       VARCHAR(64)
CATALOG_NAME          VARCHAR(64)
SCHEMA_NAME           VARCHAR(64)
TABLE_NAME            VARCHAR(64)
COLUMN_NAME           VARCHAR(64)
```

```
CURSOR_NAME           VARCHAR(64)
MESSAGE_TEXT          VARCHAR(128)
MYSQL_ERRNO           SMALLINT UNSIGNED
```

The character set for character items is UTF-8.

It is illegal to assign `NULL` to a condition information item in a `SIGNAL` statement.

A `SIGNAL` statement always specifies an `SQLSTATE` value, either directly, or indirectly by referring to a named condition defined with an `SQLSTATE` value. The first two characters of an `SQLSTATE` value are its class, and the class determines the default value for the condition information items:

- Class = `'00'` (success)

  Illegal. `SQLSTATE` values that begin with `'00'` indicate success and are not valid for `SIGNAL`.

- Class = `'01'` (warning)

  ```
  MESSAGE_TEXT = 'Unhandled user-defined warning condition';
  MYSQL_ERRNO = ER_SIGNAL_WARN
  ```

- Class = `'02'` (not found)

  ```
  MESSAGE_TEXT = 'Unhandled user-defined not found condition';
  MYSQL_ERRNO = ER_SIGNAL_NOT_FOUND
  ```

- Class > `'02'` (exception)

  ```
  MESSAGE_TEXT = 'Unhandled user-defined exception condition';
  MYSQL_ERRNO = ER_SIGNAL_EXCEPTION
  ```

For legal classes, the other condition information items are set as follows:

```
CLASS_ORIGIN = SUBCLASS_ORIGIN = '';
CONSTRAINT_CATALOG = CONSTRAINT_SCHEMA = CONSTRAINT_NAME = '';
CATALOG_NAME = SCHEMA_NAME = TABLE_NAME = COLUMN_NAME = '';
CURSOR_NAME = '';
```

The error values that are accessible after `SIGNAL` executes are the `SQLSTATE` value raised by the `SIGNAL` statement and the `MESSAGE_TEXT` and `MYSQL_ERRNO` items. These values are available from the C API:

- `SQLSTATE` value: Call `mysql_sqlstate()`
- `MYSQL_ERRNO` value: Call `mysql_errno()`
- `MESSAGE_TEXT` value: Call `mysql_error()`

From SQL, the output from `SHOW WARNINGS` and `SHOW ERRORS` indicates the `MYSQL_ERRNO` and `MESSAGE_TEXT` values in the `Code` and `Message` columns.

To retrieve information from the diagnostics area, use the `GET DIAGNOSTICS` statement (see Section 13.6.7.3, “GET DIAGNOSTICS Syntax”). For information about the diagnostics area, see Section 13.6.7.7, “The MySQL Diagnostics Area”.

Effect of Signals on Handlers, Cursors, and Statements

Signals have different effects on statement execution depending on the signal class. The class determines how severe an error is. MySQL ignores the value of the `sql_mode` system variable; in particular, strict SQL mode does not matter. MySQL also ignores `IGNORE`: The intent of `SIGNAL` is to raise a user-generated error explicitly, so a signal is never ignored.

In the following descriptions, “unhandled” means that no handler for the signaled SQLSTATE value has been defined with DECLARE ... HANDLER.

- Class = '00' (success)

  Illegal. SQLSTATE values that begin with '00' indicate success and are not valid for SIGNAL.

- Class = '01' (warning)

  The value of the warning_count system variable goes up. SHOW WARNINGS shows the signal. SQLWARNING handlers catch the signal. If the signal is unhandled in a function, statements do not end.

- Class = '02' (not found)

  NOT FOUND handlers catch the signal. There is no effect on cursors. If the signal is unhandled in a function, statements end.

- Class > '02' (exception)

  SQLEXCEPTION handlers catch the signal. If the signal is unhandled in a function, statements end.

- Class = '40'

  Treated as an ordinary exception.

Example:

```
mysql> delimiter //
mysql> CREATE FUNCTION f () RETURNS INT
    -> BEGIN
    ->   SIGNAL SQLSTATE '01234';  -- signal a warning
    ->   RETURN 5;
    -> END//
mysql> delimiter ;
mysql> CREATE TABLE t (s1 INT);
mysql> INSERT INTO t VALUES (f());
```

The result is that a row containing 5 is inserted into table t. The warning that is signaled can be viewed with SHOW WARNINGS.

13.6.7.6 Scope Rules for Handlers

A stored program may include handlers to be invoked when certain conditions occur within the program. The applicability of each handler depends on its location within the program definition and on the condition or conditions that it handles:

- A handler declared in a BEGIN ... END block is in scope only for the SQL statements following the handler declarations in the block. If the handler itself raises a condition, it cannot handle that condition, nor can any other handlers declared in the block. In the following example, handlers H1 and H2 are in scope for conditions raised by statements *stmt1* and *stmt2*. But neither H1 nor H2 are in scope for conditions raised in the body of H1 or H2.

  ```
  BEGIN -- outer block
    DECLARE EXIT HANDLER FOR ...;  -- handler H1
    DECLARE EXIT HANDLER FOR ...;  -- handler H2
    stmt1;
    stmt2;
  END;
  ```

- A handler is in scope only for the block in which it is declared, and cannot be activated for conditions occurring outside that block. In the following example, handler H1 is in scope for *stmt1* in the inner block, but not for *stmt2* in the outer block:

```
BEGIN -- outer block
  BEGIN -- inner block
    DECLARE EXIT HANDLER FOR ...;  -- handler H1
    stmt1;
  END;
  stmt2;
END;
```

- A handler can be specific or general. A specific handler is for a MySQL error code, SQLSTATE value, or condition name. A general handler is for a condition in the SQLWARNING, SQLEXCEPTION, or NOT FOUND class. Condition specificity is related to condition precedence, as described later.

Multiple handlers can be declared in different scopes and with different specificities. For example, there might be a specific MySQL error code handler in an outer block, and a general SQLWARNING handler in an inner block. Or there might be handlers for a specific MySQL error code and the general SQLWARNING class in the same block.

Whether a handler is activated depends not only on its own scope and condition value, but on what other handlers are present. When a condition occurs in a stored program, the server searches for applicable handlers in the current scope (current BEGIN ... END block). If there are no applicable handlers, the search continues outward with the handlers in each successive containing scope (block). When the server finds one or more applicable handlers at a given scope, it chooses among them based on condition precedence:

- A MySQL error code handler takes precedence over an SQLSTATE value handler.
- An SQLSTATE value handler takes precedence over general SQLWARNING, SQLEXCEPTION, or NOT FOUND handlers.
- An SQLEXCEPTION handler takes precedence over an SQLWARNING handler.
- It is possible to have several applicable handlers with the same precedence. For example, a statement could generate multiple warnings with different error codes, for each of which an error-specific handler exists. In this case, the choice of which handler the server activates is indeterminate, and may change depending on the circumstances under which the condition occurs.

One implication of the handler selection rules is that if multiple applicable handlers occur in different scopes, handlers with the most local scope take precedence over handlers in outer scopes, even over those for more specific conditions.

If there is no appropriate handler when a condition occurs, the action taken depends on the class of the condition:

- For SQLEXCEPTION conditions, the stored program terminates at the statement that raised the condition, as if there were an EXIT handler. If the program was called by another stored program, the calling program handles the condition using the handler selection rules applied to its own handlers.
- For SQLWARNING conditions, the program continues executing, as if there were a CONTINUE handler.
- For NOT FOUND conditions, if the condition was raised normally, the action is CONTINUE. If it was raised by SIGNAL or RESIGNAL, the action is EXIT.

The following examples demonstrate how MySQL applies the handler selection rules.

This procedure contains two handlers, one for the specific SQLSTATE value ('42S02') that occurs for attempts to drop a nonexistent table, and one for the general SQLEXCEPTION class:

```
CREATE PROCEDURE p1()
```

```
BEGIN
  DECLARE CONTINUE HANDLER FOR SQLSTATE '42S02'
    SELECT 'SQLSTATE handler was activated' AS msg;
  DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
    SELECT 'SQLEXCEPTION handler was activated' AS msg;

  DROP TABLE test.t;
END;
```

Both handlers are declared in the same block and have the same scope. However, `SQLSTATE` handlers take precedence over `SQLEXCEPTION` handlers, so if the table `t` is nonexistent, the `DROP TABLE` statement raises a condition that activates the `SQLSTATE` handler:

```
mysql> CALL p1();
+--------------------------------+
| msg                            |
+--------------------------------+
| SQLSTATE handler was activated |
+--------------------------------+
```

This procedure contains the same two handlers. But this time, the `DROP TABLE` statement and `SQLEXCEPTION` handler are in an inner block relative to the `SQLSTATE` handler:

```
CREATE PROCEDURE p2()
BEGIN -- outer block
    DECLARE CONTINUE HANDLER FOR SQLSTATE '42S02'
      SELECT 'SQLSTATE handler was activated' AS msg;
  BEGIN -- inner block
    DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
      SELECT 'SQLEXCEPTION handler was activated' AS msg;

    DROP TABLE test.t; -- occurs within inner block
  END;
END;
```

In this case, the handler that is more local to where the condition occurs takes precedence. The `SQLEXCEPTION` handler activates, even though it is more general than the `SQLSTATE` handler:

```
mysql> CALL p2();
+-----------------------------------+
| msg                               |
+-----------------------------------+
| SQLEXCEPTION handler was activated |
+-----------------------------------+
```

In this procedure, one of the handlers is declared in a block inner to the scope of the `DROP TABLE` statement:

```
CREATE PROCEDURE p3()
BEGIN -- outer block
  DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
    SELECT 'SQLEXCEPTION handler was activated' AS msg;
  BEGIN -- inner block
    DECLARE CONTINUE HANDLER FOR SQLSTATE '42S02'
      SELECT 'SQLSTATE handler was activated' AS msg;
  END;

  DROP TABLE test.t; -- occurs within outer block
END;
```

Only the `SQLEXCEPTION` handler applies because the other one is not in scope for the condition raised by the `DROP TABLE`:

```
mysql> CALL p3();
```

```
+-----------------------------------+
| msg                               |
+-----------------------------------+
| SQLEXCEPTION handler was activated |
+-----------------------------------+
```

In this procedure, both handlers are declared in a block inner to the scope of the `DROP TABLE` statement:

```
CREATE PROCEDURE p4()
BEGIN -- outer block
  BEGIN -- inner block
    DECLARE CONTINUE HANDLER FOR SQLEXCEPTION
      SELECT 'SQLEXCEPTION handler was activated' AS msg;
    DECLARE CONTINUE HANDLER FOR SQLSTATE '42S02'
      SELECT 'SQLSTATE handler was activated' AS msg;
  END;

  DROP TABLE test.t; -- occurs within outer block
END;
```

Neither handler applies because they are not in scope for the `DROP TABLE`. The condition raised by the statement goes unhandled and terminates the procedure with an error:

```
mysql> CALL p4();
ERROR 1051 (42S02): Unknown table 'test.t'
```

13.6.7.7 The MySQL Diagnostics Area

SQL statements produce diagnostic information that populates the diagnostics area. Standard SQL has a diagnostics area stack, containing a diagnostics area for each nested execution context. Standard SQL also supports `GET STACKED DIAGNOSTICS` syntax for referring to the second diagnostics area during condition handler execution. MySQL supports the `STACKED` keyword since MySQL 5.7.

This section describes the structure of the diagnostics area in MySQL, the information items recognized by MySQL, how statements clear and set the diagnostics area, and how diagnostics areas are pushed to and popped from the stack.

Diagnostics Area Structure

The diagnostics area contains two kinds of information:

- Statement information, such as the number of conditions that occurred or the affected-rows count.
- Condition information, such as the error code and message. If a statement raises multiple conditions, this part of the diagnostics area has a condition area for each one. If a statement raises no conditions, this part of the diagnostics area is empty.

For a statement that produces three conditions, the diagnostics area contains statement and condition information like this:

```
Statement information:
  row count
  ... other statement information items ...
Condition area list:
  Condition area 1:
    error code for condition 1
    error message for condition 1
    ... other condition information items ...
  Condition area 2:
    error code for condition 2:
    error message for condition 2
    ... other condition information items ...
```

```
Condition area 3:
  error code for condition 3
  error message for condition 3
  ... other condition information items ...
```

Diagnostics Area Information Items

The diagnostics area contains statement and condition information items. Numeric items are integers. The character set for character items is UTF-8. No item can be `NULL`. If a statement or condition item is not set by a statement that populates the diagnostics area, its value is 0 or the empty string, depending on the item data type.

The statement information part of the diagnostics area contains these items:

- `NUMBER`: An integer indicating the number of condition areas that have information.
- `ROW_COUNT`: An integer indicating the number of rows affected by the statement. `ROW_COUNT` has the same value as the `ROW_COUNT()` function (see Section 12.14, “Information Functions”).

The condition information part of the diagnostics area contains a condition area for each condition. Condition areas are numbered from 1 to the value of the `NUMBER` statement condition item. If `NUMBER` is 0, there are no condition areas.

Each condition area contains the items in the following list. All items are standard SQL except `MYSQL_ERRNO`, which is a MySQL extension. The definitions apply for conditions generated other than by a signal (that is, by a `SIGNAL` or `RESIGNAL` statement). For nonsignal conditions, MySQL populates only those condition items not described as always empty. The effects of signals on the condition area are described later.

- `CLASS_ORIGIN`: A string containing the class of the `RETURNED_SQLSTATE` value. If the `RETURNED_SQLSTATE` value begins with a class value defined in SQL standards document ISO 9075-2 (section 24.1, SQLSTATE), `CLASS_ORIGIN` is `'ISO 9075'`. Otherwise, `CLASS_ORIGIN` is `'MySQL'`.
- `SUBCLASS_ORIGIN`: A string containing the subclass of the `RETURNED_SQLSTATE` value. If `CLASS_ORIGIN` is `'ISO 9075'` or `RETURNED_SQLSTATE` ends with `'000'`, `SUBCLASS_ORIGIN` is `'ISO 9075'`. Otherwise, `SUBCLASS_ORIGIN` is `'MySQL'`.
- `RETURNED_SQLSTATE`: A string that indicates the `SQLSTATE` value for the condition.
- `MESSAGE_TEXT`: A string that indicates the error message for the condition.
- `MYSQL_ERRNO`: An integer that indicates the MySQL error code for the condition.
- `CONSTRAINT_CATALOG`, `CONSTRAINT_SCHEMA`, `CONSTRAINT_NAME`: Strings that indicate the catalog, schema, and name for a violated constraint. They are always empty.
- `CATALOG_NAME`, `SCHEMA_NAME`, `TABLE_NAME`, `COLUMN_NAME`: Strings that indicate the catalog, schema, table, and column related to the condition. They are always empty.
- `CURSOR_NAME`: A string that indicates the cursor name. This is always empty.

For the `RETURNED_SQLSTATE`, `MESSAGE_TEXT`, and `MYSQL_ERRNO` values for particular errors, see Section B.3, “Server Error Codes and Messages”.

If a `SIGNAL` (or `RESIGNAL`) statement populates the diagnostics area, its `SET` clause can assign to any condition information item except `RETURNED_SQLSTATE` any value that is legal for the item data type. `SIGNAL` also sets the `RETURNED_SQLSTATE` value, but not directly in its `SET` clause. That value comes from the `SIGNAL` statement `SQLSTATE` argument.

`SIGNAL` also sets statement information items. It sets `NUMBER` to 1. It sets `ROW_COUNT` to −1 for errors and 0 otherwise.

How the Diagnostics Area is Populated

Nondiagnostic SQL statements populate the diagnostics area automatically, and its contents can be set explicitly with the `SIGNAL` and `RESIGNAL` statements. The diagnostics area can be examined with `GET DIAGNOSTICS` to extract specific items, or with `SHOW WARNINGS` or `SHOW ERRORS` to see conditions or errors.

SQL statements clear and set the diagnostics area as follows:

- When the server starts executing a statement after parsing it, it clears the diagnostics area for nondiagnostic statements. Diagnostic statements do not clear the diagnostics area (`SHOW WARNINGS`, `SHOW ERRORS`, `GET DIAGNOSTICS`).
- If a statement raises a condition, the diagnostics area is cleared of conditions that belong to earlier statements. The exception is that conditions raised by `GET DIAGNOSTICS` and `RESIGNAL` are added to the diagnostics area without clearing it.

Thus, even a statement that does not normally clear the diagnostics area when it begins executing clears it if the statement raises a condition.

The following example shows the effect of various statements on the diagnostics area, using `SHOW WARNINGS` to display information about conditions stored there.

This `DROP TABLE` statement clears the diagnostics area and populates it when the condition occurs:

```
mysql> DROP TABLE IF EXISTS test.no_such_table;
Query OK, 0 rows affected, 1 warning (0.01 sec)

mysql> SHOW WARNINGS;
+-------+------+----------------------------------------+
| Level | Code | Message                                |
+-------+------+----------------------------------------+
| Note  | 1051 | Unknown table 'test.no_such_table' |
+-------+------+----------------------------------------+
1 row in set (0.00 sec)
```

This `SET` statement generates an error, so it clears and populates the diagnostics area:

```
mysql> SET @x = @@x;
ERROR 1193 (HY000): Unknown system variable 'x'

mysql> SHOW WARNINGS;
+-------+------+-----------------------------+
| Level | Code | Message                     |
+-------+------+-----------------------------+
| Error | 1193 | Unknown system variable 'x' |
+-------+------+-----------------------------+
1 row in set (0.00 sec)
```

The previous `SET` statement produced a single condition, so 1 is the only valid condition number for `GET DIAGNOSTICS` at this point. The following statement uses a condition number of 2, which produces a warning that is added to the diagnostics area without clearing it:

```
mysql> GET DIAGNOSTICS CONDITION 2 @p = MESSAGE_TEXT;
Query OK, 0 rows affected, 1 warning (0.00 sec)

mysql> SHOW WARNINGS;
+-------+------+-----------------------------+
| Level | Code | Message                     |
+-------+------+-----------------------------+
| Error | 1193 | Unknown system variable 'xx' |
| Error | 1753 | Invalid condition number     |
+-------+------+-----------------------------+
```

```
2 rows in set (0.00 sec)
```

Now there are two conditions in the diagnostics area, so the same `GET DIAGNOSTICS` statement succeeds:

```
mysql> GET DIAGNOSTICS CONDITION 2 @p = MESSAGE_TEXT;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @p;
+--------------------------+
| @p                       |
+--------------------------+
| Invalid condition number |
+--------------------------+
1 row in set (0.01 sec)
```

How the Diagnostics Area Stack Works

When a push to the diagnostics area stack occurs, the first (current) diagnostics area becomes the second (stacked) diagnostics area and a new current diagnostics area is created as a copy of it. Diagnostics areas are pushed to and popped from the stack under the following circumstances:

- Execution of a stored program

  A push occurs before the program executes and a pop occurs afterward. If the stored program ends while handlers are executing, there can be more than one diagnostics area to pop; this occurs due to an exception for which there are no appropriate handlers or due to `RETURN` in the handler.

  Any warning or error conditions occurring during stored program execution then are added to the current diagnostics area, except that, for triggers, only errors are added. When the stored program ends, the caller sees these conditions in its current diagonstics area.

- Execution of a condition handler within a stored program

  When a push occurs as a result of condition handler activation, the stacked diagnostics area is the area that was current within the stored program prior to the push. The new now-current diagnostics area is the handler's current diagnostics area. `GET [CURRENT] DIAGNOSTICS` and `GET STACKED DIAGNOSTICS` can be used within the handler to access the contents of the current (handler) and stacked (stored program) diagnostics areas. Initially, they return the same result, but statements executing within the handler modify the current diagnostics area, clearing and setting its contents according to the normal rules (see How the Diagnostics Area is Populated). The stacked diagnostics area cannot be modified by statements executing within the handler except `RESIGNAL`.

  If the handler executes successfully, the current (handler) diagnostics area is popped and the stacked (stored program) diagnostics area again becomes the current diagnostics area. Conditions added to the handler diagnostics area during handler execution are added to the current diagnostics area.

- Execution of `RESIGNAL`

  The `RESIGNAL` statement passes on the error condition information that is available during execution of a condition handler within a compound statement inside a stored program. `RESIGNAL` may change some or all information before passing it on, modifying the diagnostics stack as described in Section 13.6.7.4, “RESIGNAL Syntax”.

Diagnostics Area-Related System Variables

Certain system variables control or are related to some aspects of the diagnostics area:

- `max_error_count` controls the number of condition areas in the diagnostics area. If more conditions than this occur, MySQL silently discards information for the excess conditions. (Conditions

added by `RESIGNAL` are always added, with older conditions being discarded as necessary to make room.)

- `warning_count` indicates the number of conditions that occurred. This includes errors, warnings, and notes. Normally, `NUMBER` and `warning_count` are the same. However, as the number of conditions generated exceeds `max_error_count`, the value of `warning_count` continues to rise whereas `NUMBER` remains capped at `max_error_count` because no additional conditions are stored in the diagnostics area.
- `error_count` indicates the number of errors that occurred. This value includes “not found” and exception conditions, but excludes warnings and notes. Like `warning_count`, its value can exceed `max_error_count`.
- If the `sql_notes` system variable is set to 0, notes are not stored and do not increment `warning_count`.

Example: If `max_error_count` is 10, the diagnostics area can contain a maximum of 10 condition areas. Suppose that a statement raises 20 conditions, 12 of which are errors. In that case, the diagnostics area contains the first 10 conditions, `NUMBER` is 10, `warning_count` is 20, and `error_count` is 12.

Changes to the value of `max_error_count` have no effect until the next attempt to modify the diagnostics area. If the diagnostics area contains 10 condition areas and `max_error_count` is set to 5, that has no immediate effect on the size or content of the diagnostics area.

13.7 Database Administration Statements

13.7.1 Account Management Statements

MySQL account information is stored in the tables of the `mysql` database. This database and the access control system are discussed extensively in Chapter 5, *MySQL Server Administration*, which you should consult for additional details.

Important

Some MySQL releases introduce changes to the structure of the grant tables to add new privileges or features. To make sure that you can take advantage of any new capabilities, update your grant tables to have the current structure whenever you upgrade MySQL. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

When the `read_only` system variable is enabled, account-management statements require the `SUPER` privilege, in addition to any other required privileges. This is because they modify tables in the `mysql` database.

13.7.1.1 ALTER USER Syntax

```
ALTER USER [IF EXISTS]
    user [auth_option] [, user [auth_option]] ...
    [REQUIRE {NONE | tls_option [[AND] tls_option] ...}]
    [WITH resource_option [resource_option] ...]
    [password_option | lock_option] ...

ALTER USER [IF EXISTS]
    USER() IDENTIFIED BY 'auth_string'

ALTER USER [IF EXISTS]
    user DEFAULT ROLE
    {NONE | ALL | role [, role ] ...}

user:
    (see Section 6.2.4, “Specifying Account Names”)
```

```
auth_option: {
    IDENTIFIED BY 'auth_string'
  | IDENTIFIED WITH auth_plugin
  | IDENTIFIED WITH auth_plugin BY 'auth_string'
  | IDENTIFIED WITH auth_plugin AS 'hash_string'
}

tls_option: {
   SSL
 | X509
 | CIPHER 'cipher'
 | ISSUER 'issuer'
 | SUBJECT 'subject'
}

resource_option: {
    MAX_QUERIES_PER_HOUR count
  | MAX_UPDATES_PER_HOUR count
  | MAX_CONNECTIONS_PER_HOUR count
  | MAX_USER_CONNECTIONS count
}

password_option: {
    PASSWORD EXPIRE
  | PASSWORD EXPIRE DEFAULT
  | PASSWORD EXPIRE NEVER
  | PASSWORD EXPIRE INTERVAL N DAY
}

lock_option: {
    ACCOUNT LOCK
  | ACCOUNT UNLOCK
}
```

The `ALTER USER` statement modifies MySQL accounts. It enables authentication, SSL/TLS, resource-limit, and password-expiration properties to be modified for existing accounts, and enables account locking and unlocking.

Except for the `DEFAULT ROLE` syntax, `ALTER USER` requires the global `CREATE USER` privilege, or the `UPDATE` privilege for the `mysql` database. For `DEFAULT ROLE`, `ALTER USER` requires these privileges:

- To set the default roles for another user: The global `CREATE USER` privilege, or the `UPDATE` privilege for the `mysql.default_roles` system table.
- To set the default roles for yourself: No special privileges, as long as the roles you want as the default have been granted to you.

When the `read_only` system variable is enabled, `ALTER USER` additionally requires the `SUPER` privilege.

By default, an error occurs if you try to modify a user that does not exist. If the `IF EXISTS` clause is given, the statement produces a warning for each named user that does not exist, rather than an error.

The statement is written to the binary log if it succeeds, but not if it fails; in that case, rollback occurs and no changes are made. A statement written to the binary log includes all named users. If the `IF EXISTS` clause is given, this includes even users that do not exist and were not altered.

If the original statement changes the credentials for a user, the statement written to the binary log specifies the applicable authentication plugin for that user, determined as follows:

- The plugin named in the original statement, if one was specified.
- Otherwise, the plugin associated with the user account if the user exists, or the default authentication plugin if the user does not exist. (If the statement written to the binary log must specify a particular authentication plugin for a user, include it in the original statement.)

If the server adds the default authentication plugin for any users in the statement written to the binary log, it writes a warning to the error log naming those users.

`ALTER USER` modifies the `mysql.user` table row for each affected account according to the options specified in the statement. Unspecified properties retain their current values.

Example 1: Change an account's password and expire it. As a result, the user must connect with the named password and choose a new one at the next connection:

```
ALTER USER 'jeffrey'@'localhost'
  IDENTIFIED BY 'new_password' PASSWORD EXPIRE;
```

Example 2: Modify an account to use the `sha256_password` authentication plugin and the given password. Require that a new password be chosen every 180 days:

```
ALTER USER 'jeffrey'@'localhost'
  IDENTIFIED WITH sha256_password BY 'new_password'
  PASSWORD EXPIRE INTERVAL 180 DAY;
```

Example 3: Lock or unlock an account:

```
ALTER USER 'jeffrey'@'localhost' ACCOUNT LOCK;
ALTER USER 'jeffrey'@'localhost' ACCOUNT UNLOCK;
```

Example 4: Require an account to connect using SSL and establish a limit of 20 connections per hour:

```
ALTER USER 'jeffrey'@'localhost'
  REQUIRE SSL WITH MAX_CONNECTIONS_PER_HOUR 20;
```

Important

Under some circumstances, `ALTER USER` may be recorded in server logs or on the client side in a history file such as `~/.mysql_history`, which means that cleartext passwords may be read by anyone having read access to that information. For information about the conditions under which this occurs for the server logs and how to control it, see Section 6.1.2.3, “Passwords and Logging”. For similar information about client-side logging, see Section 4.5.1.3, “mysql Logging”.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. The host name part of the account name, if omitted, defaults to `'%'`. It is also possible to specify `CURRENT_USER` or `CURRENT_USER()` to refer to the account associated with the current session.

For one syntax only, the account may be specified with the `USER()` function:

```
ALTER USER USER() IDENTIFIED BY 'auth_string';
```

This syntax enables changing your own password without naming your account literally.

For `ALTER USER` syntaxes that permit an *`auth_option`* value to follow a *`user`* value, *`auth_option`* begins with `IDENTIFIED` and indicates how the account authenticates by specifying an account authentication plugin, credentials (password), or both. Each *`auth_option`* value applies *only* to the user named immediately preceding it.

Following the user specifications, the statement may include options for SSL/TLS, resource-limit, password-expiration, and locking properties. All these options are *global* to the statement and apply to *all* named users.

Example: This statement changes the password for `jeffrey` but leaves that for `jeanne` unchanged. For both accounts, connections are required to use SSL and each account can be used for a maximum of two simultaneous connections:

```
ALTER USER
  'jeffrey'@'localhost' IDENTIFIED BY 'new_password',
  'jeanne'@'localhost'
  REQUIRE SSL WITH MAX_USER_CONNECTIONS 2;
```

In the absence of a particular type of option, the account remains unchanged in that respect. For example, with no locking option, the locking state of the account is not changed.

There are several aspects to the `ALTER USER` statement, described under the following topics in this section:

- Authentication Options
- SSL/TLS Options
- Resource-Limit Options
- Password-Expiration Options
- Account-Locking Options
- Role Options

Authentication Options

An account name may be followed by an authentication option that specifies the account authentication plugin, credentials, or both:

- *`auth_plugin`* names an authentication plugin. The plugin name can be a quoted string literal or an unquoted name. Plugin names are stored in the `plugin` column of the `mysql.user` table.
- *`'auth_string'`* or *`'hash_string'`* specifiy account credentials, either as cleartext or hashed in the format expected by the authentication plugin, respectively. Credentials are stored in the `authentication_string` column of the `mysql.user` table.

`ALTER USER` permits these *`auth_option`* syntaxes:

- `IDENTIFIED BY '`*`auth_string`*`'`

  Sets the account authentication plugin to the default plugin, hashes the cleartext *`'auth_string'`* value, and stores the result in the `mysql.user` account row.

- `IDENTIFIED WITH `*`auth_plugin`*

  Sets the account authentication plugin to *`auth_plugin`*, clears the credentials to the empty string (the credentials are associated with the old authentication plugin, not the new one), and stores the result in the `mysql.user` account row.

  In addition, the password is marked expired. The user must choose a new one when next connecting.

- `IDENTIFIED WITH `*`auth_plugin`*` BY '`*`auth_string`*`'`

  Sets the account authentication plugin to *`auth_plugin`*, hashes the cleartext *`'auth_string'`* value, and stores the result in the `mysql.user` account row.

- `IDENTIFIED WITH `*`auth_plugin`*` AS '`*`hash_string`*`'`

Sets the account authentication plugin to *`auth_plugin`*, takes the hashed `'hash_string'` value as is, and stores the result in the `mysql.user` account row. The string is assumed to be already hashed in the format required by the plugin.

The default plugin is `mysql_native_password` unless the `default_authentication_plugin` system variable is set otherwise. For descriptions of each plugin, see Section 6.5.1, “Authentication Plugins”.

Example 1: Specify the password as cleartext; the default plugin is used:

```
ALTER USER 'jeffrey'@'localhost'
  IDENTIFIED BY 'mypass';
```

Example 2: Specify the authentication plugin, along with a cleartext password value:

```
ALTER USER 'jeffrey'@'localhost'
  IDENTIFIED WITH mysql_native_password
             BY 'mypass';
```

Example 3: Specify the authentication plugin, along with a hashed password value:

```
ALTER USER 'jeffrey'@'localhost'
  IDENTIFIED WITH mysql_native_password
             AS '*6C8989366EAF75BB670AD8EA7A7FC1176A95CEF4';
```

SSL/TLS Options

MySQL can check X509 certificate attributes in addition to the usual authentication that is based on the user name and credentials. For background information on the use of SSL with MySQL, see Section 6.4, “Using Secure Connections”.

To specify SSL-related options for a MySQL account, use a `REQUIRE` clause that specifies one or more *`tls_option`* values.

`ALTER USER` permits these *`tls_option`* values:

- `NONE`

  Indicates that the account has no SSL or X509 requirements. Unencrypted connections are permitted if the user name and password are valid. However, encrypted connections can also be used, at the client's option, if the client has the proper certificate and key files.

  A client need specify only the `--ssl-mode=REQUIRED` option to obtain an encrypted connection. The connection attempt fails if SSL is not available.

- `SSL`

  Tells the server to permit only encrypted connections for the account.

  ```
  ALTER USER 'jeffrey'@'localhost' REQUIRE SSL;
  ```

  A client need specify only the `--ssl-mode=REQUIRED` option to obtain an encrypted connection. The connection attempt fails if SSL is not available.

- `X509`

  Requires that the client must have a valid certificate but the exact certificate, issuer, and subject do not matter. The only requirement is that it should be possible to verify its signature with one of the CA certificates. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

```
ALTER USER 'jeffrey'@'localhost' REQUIRE X509;
```

The client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.) This is true for `ISSUER` and `SUBJECT` as well because those `REQUIRE` options imply the requirements of `X509`.

- `ISSUER 'issuer'`

  Places the restriction on connection attempts that the client must present a valid X509 certificate issued by CA `'issuer'`. If the client presents a certificate that is valid but has a different issuer, the server rejects the connection. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

  Because `ISSUER` implies the requirements of `X509`, the client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.)

```
ALTER USER 'jeffrey'@'localhost'
  REQUIRE ISSUER '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL/CN=CA/emailAddress=ca@example.com';
```

- `SUBJECT 'subject'`

  Places the restriction on connection attempts that the client must present a valid X509 certificate containing the subject `subject`. If the client presents a certificate that is valid but has a different subject, the server rejects the connection. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

  Because `SUBJECT` implies the requirements of `X509`, the client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.)

```
ALTER USER 'jeffrey'@'localhost'
  REQUIRE SUBJECT '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL demo client certificate/
    CN=client/emailAddress=client@example.com';
```

  MySQL does a simple string comparison of the `'subject'` value to the value in the certificate, so lettercase and component ordering must be given exactly as present in the certificate.

Note

Regarding `emailAddress`, see the note in the description of `REQUIRE ISSUER`.

- `CIPHER 'cipher'`

  Requests a specific cipher method for connections. This option is needed to ensure that ciphers and key lengths of sufficient strength are used. SSL itself can be weak if old algorithms using short encryption keys are used.

```
ALTER USER 'jeffrey'@'localhost' REQUIRE CIPHER 'EDH-RSA-DES-CBC3-SHA';
```

The `SUBJECT`, `ISSUER`, and `CIPHER` options can be combined in the `REQUIRE` clause:

```
ALTER USER 'jeffrey'@'localhost'
  REQUIRE SUBJECT '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL demo client certificate/
```

```
  CN=client/emailAddress=client@example.com'
AND ISSUER '/C=SE/ST=Stockholm/L=Stockholm/
  O=MySQL/CN=CA/emailAddress=ca@example.com'
AND CIPHER 'EDH-RSA-DES-CBC3-SHA';
```

The order of the options does not matter, but no option can be specified twice. The `AND` keyword is optional between `REQUIRE` options.

Resource-Limit Options

It is possible to place limits on use of server resources by an account, as discussed in Section 6.3.6, “Setting Account Resource Limits”. To do so, use a `WITH` clause that specifies one or more *`resource_option`* values:

`ALTER USER` permits these *`resource_option`* values:

- `MAX_QUERIES_PER_HOUR` *`count`*, `MAX_UPDATES_PER_HOUR` *`count`*, `MAX_CONNECTIONS_PER_HOUR` *`count`*

  These options restrict the number of queries, updates, and connections to the server permitted to this account during any given one-hour period. (Queries for which results are served from the query cache do not count against the `MAX_QUERIES_PER_HOUR` limit.) If *`count`* is `0` (the default), this means that there is no limitation for the account.

- `MAX_USER_CONNECTIONS` *`count`*

  Restricts the maximum number of simultaneous connections to the server by the account. A nonzero *`count`* specifies the limit for the account explicitly. If *`count`* is `0` (the default), the server determines the number of simultaneous connections for the account from the global value of the `max_user_connections` system variable. If `max_user_connections` is also zero, there is no limit for the account.

Example:

```
ALTER USER 'jeffrey'@'localhost'
  WITH MAX_QUERIES_PER_HOUR 500 MAX_UPDATES_PER_HOUR 100;
```

If a given resource limit is specified multiple times, the last instance takes precedence.

Password-Expiration Options

`ALTER USER` supports several *`password_option`* values for password expiration management, to either expire an account password or establish its password expiration *policy*. Policy options do not expire the password; instead, they determine how the server applies automatic expiration to the account (see Section 6.3.8, “Password Expiration Policy”).

The lifetime of a password is assessed from the date and time it was last changed.

`ALTER USER` permits these *`password_option`* values:

- `PASSWORD EXPIRE`

  Expires the account password.

  ```
  ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE;
  ```

- `PASSWORD EXPIRE DEFAULT`

  Sets the account so that the global expiration policy applies, as specified by the `default_password_lifetime` system variable.

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE DEFAULT;
```

- `PASSWORD EXPIRE NEVER`

  Disables password expiration for the account so that its password never expires.

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE NEVER;
```

- `PASSWORD EXPIRE INTERVAL N DAY`

  Sets the account password lifetime to *`N`* days. This statement requires the password to be changed every 180 days:

```
ALTER USER 'jeffrey'@'localhost' PASSWORD EXPIRE INTERVAL 180 DAY;
```

If multiple password-expiration options are specified, the last one takes precedence.

A client session operates in restricted mode if the account password was expired manually or if the password is considered past its lifetime per the automatic expiration policy. In restricted mode, operations performed within the session result in an error until the user establishes a new account password. See Section 6.3.8, “Password Expiration Policy”.

Note

It is possible to “reset” a password by setting it to its current value. As a matter of good policy, it is preferable to choose a different password.

Account-Locking Options

MySQL supports account locking and unlocking using the `ACCOUNT LOCK` and `ACCOUNT UNLOCK` options, which specify the locking state for an account. For additional discussion, see Section 6.3.12, “User Account Locking”.

If multiple account-locking options are specified, the last one takes precedence.

Role Options

`ALTER USER ... DEFAULT ROLE` defines which roles become active when the user connects to the server and authenticates, or when the user executes the `SET ROLE DEFAULT` statement during a session.

`ALTER USER ... DEFAULT ROLE` is alternative syntax for `SET DEFAULT ROLE` (see Section 13.7.1.9, “SET DEFAULT ROLE Syntax”). However, `ALTER USER` can set the default for only a single user, whereas `SET DEFAULT ROLE` can set the default for multiple users.

Each user account name uses the format described previously.

Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
ALTER USER 'joe'@'10.0.0.1' DEFAULT ROLE administrator, developer;
```

The host name part of the role name, if omitted, defaults to `'%'`.

The clause following the `DEFAULT ROLE` keywords permits these values:

- `NONE`: Set the default to `NONE` (no roles).
- `ALL`: Set the default to all roles granted to the account.

- *`role [, role ] ...`*: Set the default to the named roles, which need not exist or be granted to the account at the time `ALTER USER ... DEFAULT ROLE` is executed.

13.7.1.2 CREATE ROLE Syntax

```
CREATE ROLE [IF NOT EXISTS] role [, role ] ...
```

`CREATE ROLE` creates one or more roles, which are named collections of privileges. To use this statement, you must have the global `CREATE ROLE` or `CREATE USER` privilege.

A role when created is locked, has no password, and is assigned the default authentication plugin.

`CREATE ROLE` either succeeds for all named roles or rolls back and has no effect if any error occurs. By default, an error occurs if you try to create a role that already exists. If the `IF NOT EXISTS` clause is given, the statement produces a warning for each named role that already exists, rather than an error.

The statement is written to the binary log if it succeeds, but not if it fails; in that case, rollback occurs and no changes are made. A statement written to the binary log includes all named roles. If the `IF NOT EXISTS` clause is given, this includes even roles that already exist and were not created.

Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
CREATE ROLE 'administrator', 'developer';
CREATE ROLE 'webapp'@'localhost';
```

The host name part of the role name, if omitted, defaults to `'%'`.

For role usage examples, see Section 6.3.4, “Using Roles”.

13.7.1.3 CREATE USER Syntax

```
CREATE USER [IF NOT EXISTS]
    user [auth_option] [, user [auth_option]] ...
    [REQUIRE {NONE | tls_option [[AND] tls_option] ...}]
    [WITH resource_option [resource_option] ...]
    [password_option | lock_option] ...

user:
    (see Section 6.2.4, “Specifying Account Names”)

auth_option: {
    IDENTIFIED BY 'auth_string'
  | IDENTIFIED BY PASSWORD 'hash_string'
  | IDENTIFIED WITH auth_plugin
  | IDENTIFIED WITH auth_plugin BY 'auth_string'
  | IDENTIFIED WITH auth_plugin AS 'hash_string'
}

tls_option: {
   SSL
 | X509
 | CIPHER 'cipher'
 | ISSUER 'issuer'
 | SUBJECT 'subject'
}

resource_option: {
    MAX_QUERIES_PER_HOUR count
  | MAX_UPDATES_PER_HOUR count
  | MAX_CONNECTIONS_PER_HOUR count
  | MAX_USER_CONNECTIONS count
}
```

```
password_option: {
    PASSWORD EXPIRE
  | PASSWORD EXPIRE DEFAULT
  | PASSWORD EXPIRE NEVER
  | PASSWORD EXPIRE INTERVAL N DAY
}

lock_option: {
    ACCOUNT LOCK
  | ACCOUNT UNLOCK
}
```

The `CREATE USER` statement creates new MySQL accounts. It enables authentication, SSL/TLS, resource-limit, and password-expiration properties to be established for new accounts, and controls whether accounts are initially locked or unlocked.

An account when first created has no privileges and a default role of `NONE`.

To use `CREATE USER`, you must have the global `CREATE USER` privilege, or the `INSERT` privilege for the `mysql` database. When the `read_only` system variable is enabled, `CREATE USER` additionally requires the `SUPER` privilege.

`CREATE USER` either succeeds for all named users or rolls back and has no effect if any error occurs. By default, an error occurs if you try to create a user that already exists. If the `IF NOT EXISTS` clause is given, the statement produces a warning for each named user that already exists, rather than an error.

The statement is written to the binary log if it succeeds, but not if it fails; in that case, rollback occurs and no changes are made. A statement written to the binary log includes all named users. If the `IF NOT EXISTS` clause is given, this includes even users that already exist and were not created.

The statement written to the binary log specifies an authentication plugin for each user, determined as follows:

- The plugin named in the original statement, if one was specified.
- Otherwise, the default authentication plugin. In particular, if a user `u1` already exists and uses a nondefault authentication plugin, the statement written to the binary log for `CREATE USER IF NOT EXISTS u1` names the default authentication plugin. (If the statement written to the binary log must specify a nondefault authentication plugin for a user, include it in the original statement.)

If the server adds the default authentication plugin for any nonexisting users in the statement written to the binary log, it writes a warning to the error log naming those users.

For each account, `CREATE USER` creates a new row in the `mysql.user` table. The row reflects the properties specified in the statement. Unspecified properties are set to their default values.

Example 1: Create an account that uses the default authentication plugin and the given password. Mark the password expired so that the user must choose a new one at the first connection to the server:

```
CREATE USER 'jeffrey'@'localhost'
  IDENTIFIED BY 'new_password' PASSWORD EXPIRE;
```

Example 2: Create an account that uses the `sha256_password` authentication plugin and the given password. Require that a new password be chosen every 180 days:

```
CREATE USER 'jeffrey'@'localhost'
  IDENTIFIED WITH sha256_password BY 'new_password'
  PASSWORD EXPIRE INTERVAL 180 DAY;
```

Important

Under some circumstances, `CREATE USER` may be recorded in server logs or on the client side in a history file such as `~/.mysql_history`, which means that cleartext passwords may be read by anyone having read access to that information. For information about the conditions under which this occurs for the server logs and how to control it, see Section 6.1.2.3, “Passwords and Logging”. For similar information about client-side logging, see Section 4.5.1.3, “mysql Logging”.

For additional information about setting passwords and authentication plugins, see Section 6.3.7, “Assigning Account Passwords”, and Section 6.3.10, “Pluggable Authentication”.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. For example:

```
CREATE USER 'jeffrey'@'localhost' IDENTIFIED BY 'mypass';
```

The host name part of the account name, if omitted, defaults to `'%'`.

Each *`user`* value naming an account may be followed by an optional *`auth_option`* value that specifies how the account authenticates. These values enable account authentication plugins and credentials (passwords) to be specified. Each *`auth_option`* value applies *only* to the user named immediately preceding it.

Following the user specifications, the statement may include options for SSL/TLS, resource-limit, password-expiration, and locking properties. All these options are *global* to the statement and apply to *all* named users.

Example: This statement creates two accounts, each with the default authentication plugin and named password. For both accounts, connections must be made using a valid X509 certificate and up to 60 queries per hour are permitted. Both accounts are locked initially, so effectively they are placeholders and cannot be used until an administrator unlocks them:

```
CREATE USER
  'jeffrey'@'localhost' IDENTIFIED BY 'new_password1',
  'jeanne'@'localhost' IDENTIFIED BY 'new_password2'
  REQUIRE X509 WITH MAX_QUERIES_PER_HOUR 60
  ACCOUNT LOCK;
```

For omitted options, these default values are used:

- Authentication: The authentication plugin defined by the `default_authentication_plugin` system variable, and empty credentials
- SSL/TLS: `NONE`
- Resource limits: Unlimited
- Password expiration: `PASSWORD EXPIRE DEFAULT`
- Account locking: `ACCOUNT UNLOCK`

There are several aspects to the `CREATE USER` statement, described under the following topics in this section:

- Authentication Options
- SSL/TLS Options
- Resource-Limit Options

- Password-Expiration Options
- Account-Locking Options

Authentication Options

An account name may be followed by an authentication option that specifies the account authentication plugin, credentials, or both:

- *auth_plugin* names an authentication plugin. The plugin name can be a quoted string literal or an unquoted name. Plugin names are stored in the `plugin` column of the `mysql.user` table.
- *'auth_string'* or *'hash_string'* specifiy account credentials, either as cleartext or hashed in the format expected by the authentication plugin, respectively. Credentials are stored in the `authentication_string` column of the `mysql.user` table.

`CREATE USER` permits these *auth_option* syntaxes:

- `IDENTIFIED BY` *'auth_string'*

  Sets the account authentication plugin to the default plugin, hashes the cleartext *'auth_string'* value, and stores the result in the `mysql.user` account row.

- `IDENTIFIED BY PASSWORD` *'hash_string'*

  Sets the account authentication plugin to the default plugin, takes the hashed *'hash_string'* value as is, and stores the result in the `mysql.user` account row. The string is assumed to be already hashed in the format required by the plugin.

  Note

  This syntax is deprecated and will be removed in a future MySQL release.

- `IDENTIFIED WITH` *auth_plugin*

  Sets the account authentication plugin to *auth_plugin*, clears the credentials to the empty string, and stores the result in the `mysql.user` account row.

- `IDENTIFIED WITH` *auth_plugin* `BY` *'auth_string'*

  Sets the account authentication plugin to *auth_plugin*, hashes the cleartext *'auth_string'* value, and stores the result in the `mysql.user` account row.

- `IDENTIFIED WITH` *auth_plugin* `AS` *'hash_string'*

  Sets the account authentication plugin to *auth_plugin*, takes the *'hash_string'* value as is, and stores the result in the `mysql.user` account row. The string is assumed to be already hashed in the format required by the plugin.

The default plugin is `mysql_native_password` unless the `default_authentication_plugin` system variable is set otherwise. For descriptions of each plugin, see Section 6.5.1, “Authentication Plugins”.

Example 1: Specify the password as cleartext; the default plugin is used:

```
CREATE USER 'jeffrey'@'localhost'
  IDENTIFIED BY 'mypass';
```

Example 2: Specify the authentication plugin, along with a cleartext password value:

```
CREATE USER 'jeffrey'@'localhost'
```

```
  IDENTIFIED WITH mysql_native_password BY 'mypass';
```

SSL/TLS Options

MySQL can check X509 certificate attributes in addition to the usual authentication that is based on the user name and credentials. For background information on the use of SSL/TLS with MySQL, see Section 6.4, “Using Secure Connections”.

To specify SSL/TLS-related options for a MySQL account, use a `REQUIRE` clause that specifies one or more *`tls_option`* values:

`CREATE USER` permits these *`tls_option`* values:

- `NONE`

  Indicates that the account has no SSL or X509 requirements. Unencrypted connections are permitted if the user name and password are valid. However, encrypted connections can also be used, at the client's option, if the client has the proper certificate and key files.

  A client need specify only the `--ssl-mode=REQUIRED` option to obtain an encrypted connection. The connection attempt fails if an encrypted connection cannot be established.

- `SSL`

  Tells the server to permit only encrypted connections for the account.

  ```
  CREATE USER 'jeffrey'@'localhost' REQUIRE SSL;
  ```

  A client need specify only the `--ssl-mode=REQUIRED` option to obtain an encrypted connection. The connection attempt fails if an encrypted connection cannot be established.

- `X509`

  Requires that the client must have a valid certificate but the exact certificate, issuer, and subject do not matter. The only requirement is that it should be possible to verify its signature with one of the CA certificates. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

  ```
  CREATE USER 'jeffrey'@'localhost' REQUIRE X509;
  ```

  The client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.) This is true for `ISSUER` and `SUBJECT` as well because those `REQUIRE` options imply the requirements of `X509`.

- `ISSUER '`*`issuer`*`'`

  Places the restriction on connection attempts that the client must present a valid X509 certificate issued by CA `'`*`issuer`*`'`. If the client presents a certificate that is valid but has a different issuer, the server rejects the connection. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

  Because `ISSUER` implies the requirements of `X509`, the client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.)

  ```
  CREATE USER 'jeffrey'@'localhost'
    REQUIRE ISSUER '/C=SE/ST=Stockholm/L=Stockholm/
      O=MySQL/CN=CA/emailAddress=ca@example.com';
  ```

- `SUBJECT '`*`subject`*`'`

Places the restriction on connection attempts that the client must present a valid X509 certificate containing the subject *`subject`*. If the client presents a certificate that is valid but has a different subject, the server rejects the connection. Use of X509 certificates always implies encryption, so the `SSL` option is unnecessary in this case.

Because `SUBJECT` implies the requirements of `X509`, the client must specify the `--ssl-key` and `--ssl-cert` options to connect. (It is recommended but not required that `--ssl-ca` also be specified so that the public certificate provided by the server can be verified.)

```
CREATE USER 'jeffrey'@'localhost'
  REQUIRE SUBJECT '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL demo client certificate/
    CN=client/emailAddress=client@example.com';
```

MySQL does a simple string comparison of the *`'subject'`* value to the value in the certificate, so lettercase and component ordering must be given exactly as present in the certificate.

Note

Regarding `emailAddress`, see the note in the description of `REQUIRE ISSUER`.

- `CIPHER` *`'cipher'`*

  Requests a specific cipher method for encrypting connections. This option is needed to ensure that ciphers and key lengths of sufficient strength are used. Encryption can be weak if old algorithms using short encryption keys are used.

  ```
  CREATE USER 'jeffrey'@'localhost' REQUIRE CIPHER 'EDH-RSA-DES-CBC3-SHA';
  ```

The `SUBJECT`, `ISSUER`, and `CIPHER` options can be combined in the `REQUIRE` clause:

```
CREATE USER 'jeffrey'@'localhost'
  REQUIRE SUBJECT '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL demo client certificate/
    CN=client/emailAddress=client@example.com'
  AND ISSUER '/C=SE/ST=Stockholm/L=Stockholm/
    O=MySQL/CN=CA/emailAddress=ca@example.com'
  AND CIPHER 'EDH-RSA-DES-CBC3-SHA';
```

The order of the options does not matter, but no option can be specified twice. The `AND` keyword is optional between `REQUIRE` options.

Resource-Limit Options

It is possible to place limits on use of server resources by an account, as discussed in Section 6.3.6, “Setting Account Resource Limits”. To do so, use a `WITH` clause that specifies one or more *`resource_option`* values.

`CREATE USER` permits these *`resource_option`* values:

- `MAX_QUERIES_PER_HOUR` *`count`*, `MAX_UPDATES_PER_HOUR` *`count`*, `MAX_CONNECTIONS_PER_HOUR` *`count`*

  These options restrict the number of queries, updates, and connections to the server permitted to this account during any given one-hour period. (Queries for which results are served from the query cache do not count against the `MAX_QUERIES_PER_HOUR` limit.) If *`count`* is `0` (the default), this means that there is no limitation for the account.

- `MAX_USER_CONNECTIONS` *`count`*

Restricts the maximum number of simultaneous connections to the server by the account. A nonzero *`count`* specifies the limit for the account explicitly. If *`count`* is `0` (the default), the server determines the number of simultaneous connections for the account from the global value of the `max_user_connections` system variable. If `max_user_connections` is also zero, there is no limit for the account.

Example:

```
CREATE USER 'jeffrey'@'localhost'
  WITH MAX_QUERIES_PER_HOUR 500 MAX_UPDATES_PER_HOUR 100;
```

If a given resource limit is specified multiple times, the last instance takes precedence.

Password-Expiration Options

`CREATE USER` supports several *`password_option`* values for password expiration management, to either expire an account password or establish its password expiration *policy*. Policy options do not expire the password; instead, they determine how the server applies automatic expiration to the account (see Section 6.3.8, “Password Expiration Policy”).

The lifetime of a password is assessed from the date and time it was last changed.

`CREATE USER` permits these *`password_option`* values:

- `PASSWORD EXPIRE`

  Expires the account password.

  ```
  CREATE USER 'jeffrey'@'localhost' PASSWORD EXPIRE;
  ```

- `PASSWORD EXPIRE DEFAULT`

  Sets the account so that the global expiration policy applies, as specified by the `default_password_lifetime` system variable.

  ```
  CREATE USER 'jeffrey'@'localhost' PASSWORD EXPIRE DEFAULT;
  ```

- `PASSWORD EXPIRE NEVER`

  Disables password expiration for the account so that its password never expires.

  ```
  CREATE USER 'jeffrey'@'localhost' PASSWORD EXPIRE NEVER;
  ```

- `PASSWORD EXPIRE INTERVAL` *`N`* `DAY`

  Sets the account password lifetime to *`N`* days. This statement requires the password to be changed every 180 days:

  ```
  CREATE USER 'jeffrey'@'localhost' PASSWORD EXPIRE INTERVAL 180 DAY;
  ```

If multiple password-expiration options are specified, the last one takes precedence.

A client session operates in restricted mode if the account password was expired manually or if the password is considered past its lifetime per the automatic expiration policy. In restricted mode, operations performed within the session result in an error until the user establishes a new account password. See Section 6.3.8, “Password Expiration Policy”.

Account-Locking Options

MySQL supports account locking and unlocking using the `ACCOUNT LOCK` and `ACCOUNT UNLOCK` options, which specify the locking state for an account. For additional discussion, see Section 6.3.12, “User Account Locking”.

If multiple account-locking options are specified, the last one takes precedence.

13.7.1.4 DROP ROLE Syntax

```
DROP ROLE [IF EXISTS] role [, role ] ...
```

`DROP ROLE` removes one or more roles (named collections of privileges). To use this statement, you must have the global `DROP ROLE` or `CREATE USER` privilege.

`DROP ROLE` either succeeds for all named roles or rolls back and has no effect if any error occurs. By default, an error occurs if you try to drop a role that does not exist. If the `IF EXISTS` clause is given, the statement produces a warning for each named role that does not exist, rather than an error.

The statement is written to the binary log if it succeeds, but not if it fails; in that case, rollback occurs and no changes are made. A statement written to the binary log includes all named roles. If the `IF EXISTS` clause is given, this includes even roles that do not exist and were not dropped.

Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
DROP ROLE 'administrator', 'developer';
DROP ROLE 'webapp'@'localhost';
```

The host name part of the role name, if omitted, defaults to `'%'`.

A dropped role is automatically revoked from any user account (or role) to which the role was granted. Within any current session for such an account, its privileges are adjusted for the next statement executed.

13.7.1.5 DROP USER Syntax

```
DROP USER [IF EXISTS] user [, user] ...
```

The `DROP USER` statement removes one or more MySQL accounts and their privileges. It removes privilege rows for the account from all grant tables.

To use `DROP USER`, you must have the global `CREATE USER` privilege, or the `DELETE` privilege for the `mysql` database. When the `read_only` system variable is enabled, `DROP USER` additionally requires the `SUPER` privilege.

`DROP USER` either succeeds for all named users or rolls back and has no effect if any error occurs. By default, an error occurs if you try to drop a user that does not exist. If the `IF EXISTS` clause is given, the statement produces a warning for each named user that does not exist, rather than an error.

The statement is written to the binary log if it succeeds, but not if it fails; in that case, rollback occurs and no changes are made. A statement written to the binary log includes all named users. If the `IF EXISTS` clause is given, this includes even users that do not exist and were not dropped.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. For example:

```
DROP USER 'jeffrey'@'localhost';
```

The host name part of the account name, if omitted, defaults to `'%'`.

Important

`DROP USER` does not automatically close any open user sessions. Rather, in the event that a user with an open session is dropped, the statement does not take effect until that user's session is closed. Once the session is closed, the user is dropped, and that user's next attempt to log in will fail. *This is by design*.

`DROP USER` does not automatically drop or invalidate databases or objects within them that the old user created. This includes stored programs or views for which the `DEFINER` attribute names the dropped user. Attempts to access such objects may produce an error if they execute in definer security context. (For information about security context, see Section 23.6, “Access Control for Stored Programs and Views”.)

13.7.1.6 GRANT Syntax

```
GRANT
    priv_type [(column_list)]
      [, priv_type [(column_list)]] ...
    ON [object_type] priv_level
    TO user_or_role [auth_option] [, user_or_role [auth_option]] ...
    [REQUIRE {NONE | tls_option [[AND] tls_option] ...}]
    [WITH {GRANT OPTION | resource_option} ...]

GRANT PROXY ON user_or_role
    TO user_or_role [, user_or_role] ...
    [WITH GRANT OPTION]

GRANT role [, role] ...
    TO user_or_role [, user_or_role] ...
    [WITH ADMIN OPTION]

object_type: {
    TABLE
  | FUNCTION
  | PROCEDURE
}

priv_level: {
    *
  | *.*
  | db_name.*
  | db_name.tbl_name
  | tbl_name
  | db_name.routine_name
}

user_or_role: {
    user
  | role
}

user:
    (see Section 6.2.4, “Specifying Account Names”)

role:
    (see Section 6.2.5, “Specifying Role Names”.

auth_option: {
    IDENTIFIED BY 'auth_string'
  | IDENTIFIED BY PASSWORD 'hash_string'
  | IDENTIFIED WITH auth_plugin
  | IDENTIFIED WITH auth_plugin BY 'auth_string'
  | IDENTIFIED WITH auth_plugin AS 'hash_string'
}

tls_option: {
    SSL
  | X509
  | CIPHER 'cipher'
```

```
  | ISSUER 'issuer'
  | SUBJECT 'subject'
}

resource_option: {
  | MAX_QUERIES_PER_HOUR count
  | MAX_UPDATES_PER_HOUR count
  | MAX_CONNECTIONS_PER_HOUR count
  | MAX_USER_CONNECTIONS count
}
```

The GRANT statement enables system administrators to grant privileges and roles, which can be granted to user accounts and roles. These syntax restrictions apply:

- GRANT cannot mix granting both privileges and roles in the same statement. A given GRANT statement must grant either privileges or roles.
- The ON clause distinguishes whether the statement grants privileges or roles:
  - With ON, the statement grants privileges.
  - Without ON, the statement grants roles.
  - It is permitted to assign both privileges and roles to an account, but you must use separate GRANT statements, each with syntax appropriate to what is to be granted.

For more information about roles, see Section 6.3.4, “Using Roles”.

To use GRANT, you must have the GRANT OPTION privilege, and you must have the privileges that you are granting. When the read_only system variable is enabled, GRANT additionally requires the SUPER privilege.

GRANT either succeeds for all named users and roles or rolls back and has no effect if any error occurs. The statement is written to the binary log only if it succeeds for all named users and roles.

The REVOKE statement is related to GRANT and enables administrators to remove account privileges. See Section 13.7.1.8, “REVOKE Syntax”.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
GRANT ALL ON db1.* TO 'jeffrey'@'localhost';
GRANT 'role1', 'role2' TO 'user1'@'localhost', 'user2'@'localhost';
GRANT SELECT ON world.* TO 'role3';
```

The host name part of the account or role name, if omitted, defaults to '%'.

Normally, a database administrator first uses CREATE USER to create an account and define its nonprivilege characteristics such as its password, whether it uses secure connections, and limits on access to server resources, then uses GRANT to define its privileges. ALTER USER may be used to change the nonprivilege characteristics of existing accounts. For example:

```
CREATE USER 'jeffrey'@'localhost' IDENTIFIED BY 'mypass';
GRANT ALL ON db1.* TO 'jeffrey'@'localhost';
GRANT SELECT ON db2.invoice TO 'jeffrey'@'localhost';
ALTER USER 'jeffrey'@'localhost' WITH MAX_QUERIES_PER_HOUR 90;
```

Note

Examples shown here include no IDENTIFIED clause. It is assumed that you establish passwords with CREATE USER at account-creation time to avoid creating insecure accounts.

Note

If an account named in a `GRANT` statement does not already exist, `GRANT` may create it under the conditions described later in the discussion of the `NO_AUTO_CREATE_USER` SQL mode. It is also possible to use `GRANT` to specify nonprivilege account characteristics such as whether it uses secure connections and limits on access to server resources.

However, use of `GRANT` to create accounts or define nonprivilege characteristics is deprecated. Instead, perform these tasks using `CREATE USER` or `ALTER USER`.

From the `mysql` program, `GRANT` responds with `Query OK, 0 rows affected` when executed successfully. To determine what privileges result from the operation, use `SHOW GRANTS`. See Section 13.7.5.21, “SHOW GRANTS Syntax”.

There are several aspects to the `GRANT` statement, described under the following topics in this section:

- Object Quoting Guidelines
- Privileges Supported by MySQL
- Global Privileges
- Database Privileges
- Table Privileges
- Column Privileges
- Stored Routine Privileges
- Proxy User Privileges
- Granting Roles
- Account Names and Passwords
- Implicit Account Creation
- Other Account Characteristics
- MySQL and Standard SQL Versions of GRANT

`GRANT` supports host names up to 60 characters long. User names can be up to 32 characters. Database, table, column, and routine names can be up to 64 characters.

Warning

Do not attemt to change the permissible length for user names by altering the `mysql.user` table. Doing so results in unpredictable behavior which may even make it impossible for users to log in to the MySQL server. Never alter the structure of tables in the `mysql` database in any manner except by means of the procedure described in Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

Important

Under some circumstances, `GRANT` may be recorded in server logs or on the client side in a history file such as `~/.mysql_history`, which means that cleartext passwords may be read by anyone having read access to that information. For information about the conditions under which this occurs for the

server logs and how to control it, see Section 6.1.2.3, “Passwords and Logging”. For similar information about client-side logging, see Section 4.5.1.3, “mysql Logging”.

Object Quoting Guidelines

Several objects within `GRANT` statements are subject to quoting, although quoting is optional in many cases: Account, role, database, table, column, and routine names. For example, if a *`user_name`* or *`host_name`* value in an account name is legal as an unquoted identifier, you need not quote it. However, quotation marks are necessary to specify a *`user_name`* string containing special characters (such as `-`), or a *`host_name`* string containing special characters or wildcard characters (such as `%`); for example, `'test-user'@'%.com'`. Quote the user name and host name separately.

To specify quoted values, quote database, table, column, and routine names as identifiers. Quote user names and host names as identifiers or as strings. Quote passwords as strings. For string-quoting and identifier-quoting guidelines, see Section 9.1.1, “String Literals”, and Section 9.2, “Schema Object Names”.

The `_` and `%` wildcards are permitted when specifying database names in `GRANT` statements that grant privileges at the database level. This means, for example, that if you want to use a `_` character as part of a database name, you should specify it as `\_` in the `GRANT` statement, to prevent the user from being able to access additional databases matching the wildcard pattern; for example, `GRANT ... ON \`foo_bar\`.* TO ...`.

Privileges Supported by MySQL

The following tables summarize the permissible static and dynamic *`priv_type`* privilege types that can be specified for the `GRANT` and `REVOKE` statements, and the levels at which each privilege can be granted. For additional information about each privilege, see Section 6.2.1, “Privileges Provided by MySQL”. For information about the differences between static and dynamic privileges, see Section 6.2.2, “Static Versus Dynamic Privileges”.

Table 13.3 Permissible Static Privileges for GRANT and REVOKE

Privilege	Meaning and Grantable Levels
`ALL [PRIVILEGES]`	Grant all privileges at specified access level except `GRANT OPTION` and `PROXY`.
`ALTER`	Enable use of `ALTER TABLE`. Levels: Global, database, table.
`ALTER ROUTINE`	Enable stored routines to be altered or dropped. Levels: Global, database, procedure.
`CREATE`	Enable database and table creation. Levels: Global, database, table.
`CREATE ROUTINE`	Enable stored routine creation. Levels: Global, database.
`CREATE TABLESPACE`	Enable tablespaces and log file groups to be created, altered, or dropped. Level: Global.
`CREATE TEMPORARY TABLES`	Enable use of `CREATE TEMPORARY TABLE`. Levels: Global, database.
`CREATE USER`	Enable use of `CREATE USER`, `DROP USER`, `RENAME USER`, and `REVOKE ALL PRIVILEGES`. Level: Global.
`CREATE VIEW`	Enable views to be created or altered. Levels: Global, database, table.
`DELETE`	Enable use of `DELETE`. Level: Global, database, table.
`DROP`	Enable databases, tables, and views to be dropped. Levels: Global, database, table.
`EVENT`	Enable use of events for the Event Scheduler. Levels: Global, database.

Privilege	Meaning and Grantable Levels
EXECUTE	Enable the user to execute stored routines. Levels: Global, database, table.
FILE	Enable the user to cause the server to read or write files. Level: Global.
GRANT OPTION	Enable privileges to be granted to or removed from other accounts. Levels: Global, database, table, procedure, proxy.
INDEX	Enable indexes to be created or dropped. Levels: Global, database, table.
INSERT	Enable use of INSERT. Levels: Global, database, table, column.
LOCK TABLES	Enable use of LOCK TABLES on tables for which you have the SELECT privilege. Levels: Global, database.
PROCESS	Enable the user to see all processes with SHOW PROCESSLIST. Level: Global.
PROXY	Enable user proxying. Level: From user to user.
REFERENCES	Enable foreign key creation. Levels: Global, database, table, column.
RELOAD	Enable use of FLUSH operations. Level: Global.
REPLICATION CLIENT	Enable the user to ask where master or slave servers are. Level: Global.
REPLICATION SLAVE	Enable replication slaves to read binary log events from the master. Level: Global.
SELECT	Enable use of SELECT. Levels: Global, database, table, column.
SHOW DATABASES	Enable SHOW DATABASES to show all databases. Level: Global.
SHOW VIEW	Enable use of SHOW CREATE VIEW. Levels: Global, database, table.
SHUTDOWN	Enable use of mysqladmin shutdown. Level: Global.
SUPER	Enable use of other administrative operations such as CHANGE MASTER TO, KILL, PURGE BINARY LOGS, SET GLOBAL, and mysqladmin debug command. Level: Global.
TRIGGER	Enable trigger operations. Levels: Global, database, table.
UPDATE	Enable use of UPDATE. Levels: Global, database, table, column.
USAGE	Synonym for “no privileges”

Table 13.4 Permissible Dynamic Privileges for GRANT and REVOKE

Privilege	Meaning and Grantable Levels
AUDIT_ADMIN	Enable audit log configuration. Level: Global.
BINLOG_ADMIN	Enable binary log control. Level: Global.
CONNECTION_ADMIN	Enable connection limit/restriction control. Level: Global.
ENCRYPTION_KEY_ADMIN	Enable InnoDB key rotation. Level: Global.
FIREWALL_ADMIN	Enable firewall rule administration, any user. Level: Global.
FIREWALL_USER	Enable firewall rule administration, self. Level: Global.
GROUP_REPLICATION_ADMIN	Enable Group Replication control. Level: Global.
REPLICATION_SLAVE_ADMIN	Enable regular replication control. Level: Global.
ROLE_ADMIN	Enable use of WITH ADMIN OPTION. Level: Global.
SET_USER_ID	Enable setting non-self DEFINER values. Level: Global.
SYSTEM_VARIABLES_ADMIN	Enable modifying or persisting global system variables. Level: Global.
VERSION_TOKEN_ADMIN	Enable use of Version Tokens UDFs. Level: Global.

A trigger is associated with a table. To create or drop a trigger, you must have the `TRIGGER` privilege for the table, not the trigger.

In `GRANT` statements, the `ALL [PRIVILEGES]` or `PROXY` privilege must be named by itself and cannot be specified along with other privileges. `ALL [PRIVILEGES]` stands for all privileges available for the level at which privileges are to be granted except for the `GRANT OPTION` and `PROXY` privileges.

`USAGE` can be specified to create a user that has no privileges, or to specify the `REQUIRE` or `WITH` clauses for an account without changing its existing privileges. (However, use of `GRANT` to define nonprivilege characteristics is deprecated.

MySQL account information is stored in the tables of the `mysql` database. For additional details, consult Section 6.2, “The MySQL Access Privilege System”, which discusses the `mysql` database and the access control system extensively.

If the grant tables hold privilege rows that contain mixed-case database or table names and the `lower_case_table_names` system variable is set to a nonzero value, `REVOKE` cannot be used to revoke these privileges. It will be necessary to manipulate the grant tables directly. (`GRANT` will not create such rows when `lower_case_table_names` is set, but such rows might have been created prior to setting that variable.)

Privileges can be granted at several levels, depending on the syntax used for the `ON` clause. For `REVOKE`, the same `ON` syntax specifies which privileges to remove.

For the global, database, table, and routine levels, `GRANT ALL` assigns only the privileges that exist at the level you are granting. For example, `GRANT ALL ON db_name.*` is a database-level statement, so it does not grant any global-only privileges such as `FILE`. Granting `ALL` does not assign the `GRANT OPTION` or `PROXY` privilege.

The *`object_type`* clause, if present, should be specified as `TABLE`, `FUNCTION`, or `PROCEDURE` when the following object is a table, a stored function, or a stored procedure.

The privileges for a database, table, column, or routine are formed additively as the logical `OR` of the privileges at each of the privilege levels. For example, if a user has a global `SELECT` privilege, the privilege cannot be denied by an absence of the privilege at the database, table, or column level. Details of the privilege-checking procedure are presented in Section 6.2.7, “Access Control, Stage 2: Request Verification”.

If you are using table, column, or routine privileges for even one user, the server examines table, column, and routine privileges for all users and this slows down MySQL a bit. Similarly, if you limit the number of queries, updates, or connections for any users, the server must monitor these values.

MySQL enables you to grant privileges on databases or tables that do not exist. For tables, the privileges to be granted must include the `CREATE` privilege. *This behavior is by design*, and is intended to enable the database administrator to prepare user accounts and privileges for databases or tables that are to be created at a later time.

Important

MySQL does not automatically revoke any privileges when you drop a database or table. However, if you drop a routine, any routine-level privileges granted for that routine are revoked.

Global Privileges

Global privileges are administrative or apply to all databases on a given server. To assign global privileges, use `ON *.*` syntax:

```
GRANT ALL ON *.* TO 'someuser'@'somehost';
GRANT SELECT, INSERT ON *.* TO 'someuser'@'somehost';
```

The CREATE TABLESPACE, CREATE USER, FILE, PROCESS, RELOAD, REPLICATION CLIENT, REPLICATION SLAVE, SHOW DATABASES, SHUTDOWN, and SUPER static privileges are administrative and can only be granted globally.

Dynamic privileges are all global and can only be granted globally.

Other privileges can be granted globally or at more specific levels.

The effect of GRANT OPTION granted at the global level differs for static and dynamic privileges:

- GRANT OPTION granted for any static global privilege applies to all static global privileges.
- GRANT OPTION granted for any dynamic privilege applies only to that dynamic privilege.

GRANT ALL at the global level grants all static global privileges and all currently registered dynamic privileges. A dynamic privilege registered subsequent to execution of the GRANT statement is not granted retroactively to any account.

MySQL stores global privileges in the mysql.user table.

Database Privileges

Database privileges apply to all objects in a given database. To assign database-level privileges, use ON *db_name.** syntax:

```
GRANT ALL ON mydb.* TO 'someuser'@'somehost';
GRANT SELECT, INSERT ON mydb.* TO 'someuser'@'somehost';
```

If you use ON * syntax (rather than ON *.*) and you have selected a default database, privileges are assigned at the database level for the default database. An error occurs if there is no default database.

The CREATE, DROP, EVENT, GRANT OPTION, LOCK TABLES, and REFERENCES privileges can be specified at the database level. Table or routine privileges also can be specified at the database level, in which case they apply to all tables or routines in the database.

MySQL stores database privileges in the mysql.db table.

Table Privileges

Table privileges apply to all columns in a given table. To assign table-level privileges, use ON *db_name.tbl_name* syntax:

```
GRANT ALL ON mydb.mytbl TO 'someuser'@'somehost';
GRANT SELECT, INSERT ON mydb.mytbl TO 'someuser'@'somehost';
```

If you specify *tbl_name* rather than *db_name.tbl_name*, the statement applies to *tbl_name* in the default database. An error occurs if there is no default database.

The permissible *priv_type* values at the table level are ALTER, CREATE VIEW, CREATE, DELETE, DROP, GRANT OPTION, INDEX, INSERT, REFERENCES, SELECT, SHOW VIEW, TRIGGER, and UPDATE.

Table-level privileges apply to base tables and views. They do not apply to tables created with CREATE TEMPORARY TABLE, even if the table names match. For information about TEMPORARY table privileges, see Section 13.1.15.3, “CREATE TEMPORARY TABLE Syntax”.

MySQL stores table privileges in the mysql.tables_priv table.

Column Privileges

Column privileges apply to single columns in a given table. Each privilege to be granted at the column level must be followed by the column or columns, enclosed within parentheses.

```
GRANT SELECT (col1), INSERT (col1,col2) ON mydb.mytbl TO 'someuser'@'somehost';
```

The permissible *`priv_type`* values for a column (that is, when you use a *`column_list`* clause) are `INSERT`, `REFERENCES`, `SELECT`, and `UPDATE`.

MySQL stores column privileges in the `mysql.columns_priv` table.

Stored Routine Privileges

The `ALTER ROUTINE`, `CREATE ROUTINE`, `EXECUTE`, and `GRANT OPTION` privileges apply to stored routines (procedures and functions). They can be granted at the global and database levels. Except for `CREATE ROUTINE`, these privileges can be granted at the routine level for individual routines.

```
GRANT CREATE ROUTINE ON mydb.* TO 'someuser'@'somehost';
GRANT EXECUTE ON PROCEDURE mydb.myproc TO 'someuser'@'somehost';
```

The permissible *`priv_type`* values at the routine level are `ALTER ROUTINE`, `EXECUTE`, and `GRANT OPTION`. `CREATE ROUTINE` is not a routine-level privilege because you must have this privilege to create a routine in the first place.

MySQL stores routine-level privileges in the `mysql.procs_priv` table.

Proxy User Privileges

The `PROXY` privilege enables one user to be a proxy for another. The proxy user impersonates or takes the identity of the proxied user.

```
GRANT PROXY ON 'localuser'@'localhost' TO 'externaluser'@'somehost';
```

When `PROXY` is granted, it must be the only privilege named in the `GRANT` statement, the `REQUIRE` clause cannot be given, and the only permitted `WITH` option is `WITH GRANT OPTION`.

Proxying requires that the proxy user authenticate through a plugin that returns the name of the proxied user to the server when the proxy user connects, and that the proxy user have the `PROXY` privilege for the proxied user. For details and examples, see Section 6.3.11, “Proxy Users”.

MySQL stores proxy privileges in the `mysql.proxies_priv` table.

Granting Roles

`GRANT` syntax without an `ON` clause grants roles rather than individual privileges. A role is a named collection of privileges; see Section 6.3.4, “Using Roles”. For example:

```
GRANT 'role1', 'role2' TO 'user1'@'localhost', 'user2'@'localhost';
```

Each role to be granted must exist, as well as each user account or role to which it is to be granted.

If the `GRANT` statement includes the `WITH ADMIN OPTION` clause, each named user becomes able to grant the named roles to other users or roles, or revoke them from other users or roles. This includes the ability to use `WITH ADMIN OPTION` itself.

It is possible to create circular references with `GRANT`. For example:

```
CREATE USER 'u1', 'u2';
CREATE ROLE 'r1', 'r2';

GRANT 'u1' TO 'u1';   -- simple loop: u1 -> u1
GRANT 'r1' TO 'r1';   -- simple loop: r1 -> r1
```

```
GRANT 'r2' TO 'u2';
GRANT 'u2' TO 'r2';   -- mixed user/role loop: u2 -> r2 -> u2
```

Circular grant references are permitted but add no new privileges or roles to the grantee because a user or role already has its privileges and roles.

Account Names and Passwords

A *`user`* value in a `GRANT` statement indicates a MySQL account to which the statement applies. To accommodate granting rights to users from arbitrary hosts, MySQL supports specifying the *`user`* value in the form *`'user_name'@'host_name'`*.

You can specify wildcards in the host name. For example, *`'user_name'@'%.example.com'`* applies to *`user_name`* for any host in the `example.com` domain, and *`'user_name'@'192.168.1.%'`* applies to *`user_name`* for any host in the `192.168.1` class C subnet.

The simple form *`'user_name'`* is a synonym for *`'user_name'@'%'`*.

MySQL does not support wildcards in user names. To refer to an anonymous user, specify an account with an empty user name with the `GRANT` statement:

```
GRANT ALL ON test.* TO ''@'localhost' ...;
```

In this case, any user who connects from the local host with the correct password for the anonymous user will be permitted access, with the privileges associated with the anonymous-user account.

For additional information about user name and host name values in account names, see Section 6.2.4, “Specifying Account Names”.

Warning

If you permit local anonymous users to connect to the MySQL server, you should also grant privileges to all local users as *`'user_name'@'localhost'`*. Otherwise, the anonymous user account for `localhost` in the `mysql.user` table is used when named users try to log in to the MySQL server from the local machine. For details, see Section 6.2.6, “Access Control, Stage 1: Connection Verification”.

To determine whether this issue applies to you, execute the following query, which lists any anonymous users:

```
SELECT Host, User FROM mysql.user WHERE User='';
```

To avoid the problem just described, delete the local anonymous user account using this statement:

```
DROP USER ''@'localhost';
```

For `GRANT` syntaxes that permit an *`auth_option`* value to follow a *`user`* value, *`auth_option`* begins with `IDENTIFIED` and indicates how the account authenticates by specifying an account authentication plugin, credentials (password), or both. Syntax of the *`auth_option`* clause is the same as for the `CREATE USER` statement. For details, see Section 13.7.1.3, “CREATE USER Syntax”.

Note

Use of `GRANT` to define account authentication characteristics is deprecated. Instead, establish or change authentication characteristics using `CREATE USER`

or `ALTER USER`. This `GRANT` capability will be removed in a future MySQL release.

When `IDENTIFIED` is present and you have the global grant privilege (`GRANT OPTION`), any password specified becomes the new password for the account, even if the account exists and already has a password. Without `IDENTIFIED`, the account password remains unchanged.

Implicit Account Creation

If an account named in a `GRANT` statement does not exist, the action taken depends on the `NO_AUTO_CREATE_USER` SQL mode:

- If `NO_AUTO_CREATE_USER` is not enabled, `GRANT` creates the account. *This is very insecure* unless you specify a nonempty password using `IDENTIFIED BY`.
- If `NO_AUTO_CREATE_USER` is enabled, `GRANT` fails and does not create the account, unless you specify a nonempty password using `IDENTIFIED BY` or name an authentication plugin using `IDENTIFIED WITH`.

If the account already exists, `IDENTIFIED WITH` is prohibited because it is intended only for use when creating new accounts.

Other Account Characteristics

MySQL can check X509 certificate attributes in addition to the usual authentication that is based on the user name and credentials. For background information on the use of SSL with MySQL, see Section 6.4, “Using Secure Connections”.

The optional `REQUIRE` clause specifies SSL-related options for a MySQL account. The syntax is the same as for the `CREATE USER` statement. For details, see Section 13.7.1.3, “CREATE USER Syntax”.

Note

Use of `GRANT` to define account SSL characteristics is deprecated. Instead, establish or change SSL characteristics using `CREATE USER` or `ALTER USER`. This `GRANT` capability will be removed in a future MySQL release.

The optional `WITH` clause is used for these purposes:

- To enable a user to grant privileges to other users
- To specify resource limits for a user

The `WITH GRANT OPTION` clause gives the user the ability to give to other users any privileges the user has at the specified privilege level.

To grant the `GRANT OPTION` privilege to an account without otherwise changing its privileges, do this:

```
GRANT USAGE ON *.* TO 'someuser'@'somehost' WITH GRANT OPTION;
```

Be careful to whom you give the `GRANT OPTION` privilege because two users with different privileges may be able to combine privileges!

You cannot grant another user a privilege which you yourself do not have; the `GRANT OPTION` privilege enables you to assign only those privileges which you yourself possess.

Be aware that when you grant a user the `GRANT OPTION` privilege at a particular privilege level, any privileges the user possesses (or may be given in the future) at that level can also be granted by that user to other users. Suppose that you grant a user the `INSERT` privilege on a database. If you then

grant the SELECT privilege on the database and specify WITH GRANT OPTION, that user can give to other users not only the SELECT privilege, but also INSERT. If you then grant the UPDATE privilege to the user on the database, the user can grant INSERT, SELECT, and UPDATE.

For a nonadministrative user, you should not grant the ALTER privilege globally or for the mysql database. If you do that, the user can try to subvert the privilege system by renaming tables!

For additional information about security risks associated with particular privileges, see Section 6.2.1, “Privileges Provided by MySQL”.

It is possible to place limits on use of server resources by an account, as discussed in Section 6.3.6, “Setting Account Resource Limits”. To do so, use a WITH clause that specifies one or more *resource_option* values. Limits not specified retain their current values. The syntax is the same as for the CREATE USER statement. For details, see Section 13.7.1.3, “CREATE USER Syntax”.

Note

Use of GRANT to define account resource limits is deprecated. Instead, establish or change resource limits using CREATE USER or ALTER USER. This GRANT capability will be removed in a future MySQL release.

MySQL and Standard SQL Versions of GRANT

The biggest differences between the MySQL and standard SQL versions of GRANT are:

- MySQL associates privileges with the combination of a host name and user name and not with only a user name.
- Standard SQL does not have global or database-level privileges, nor does it support all the privilege types that MySQL supports.
- MySQL does not support the standard SQL UNDER privilege.
- Standard SQL privileges are structured in a hierarchical manner. If you remove a user, all privileges the user has been granted are revoked. This is also true in MySQL if you use DROP USER. See Section 13.7.1.5, “DROP USER Syntax”.
- In standard SQL, when you drop a table, all privileges for the table are revoked. In standard SQL, when you revoke a privilege, all privileges that were granted based on that privilege are also revoked. In MySQL, privileges can be dropped with DROP USER or REVOKE statements.
- In MySQL, it is possible to have the INSERT privilege for only some of the columns in a table. In this case, you can still execute INSERT statements on the table, provided that you insert values only for those columns for which you have the INSERT privilege. The omitted columns are set to their implicit default values if strict SQL mode is not enabled. In strict mode, the statement is rejected if any of the omitted columns have no default value. (Standard SQL requires you to have the INSERT privilege on all columns.) For information about strict SQL mode and implicit default values, see Section 5.1.8, “Server SQL Modes”, and Section 11.7, “Data Type Default Values”.

13.7.1.7 RENAME USER Syntax

```
RENAME USER old_user TO new_user
    [, old_user TO new_user] ...
```

The RENAME USER statement renames existing MySQL accounts. An error occurs for old accounts that do not exist or new accounts that already exist.

To use RENAME USER, you must have the global CREATE USER privilege, or the UPDATE privilege for the mysql database. When the read_only system variable is enabled, RENAME USER additionally requires the SUPER privilege.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. For example:

```
RENAME USER 'jeffrey'@'localhost' TO 'jeff'@'127.0.0.1';
```

The host name part of the account name, if omitted, defaults to `'%'`.

`RENAME USER` causes the privileges held by the old user to be those held by the new user. However, `RENAME USER` does not automatically drop or invalidate databases or objects within them that the old user created. This includes stored programs or views for which the `DEFINER` attribute names the old user. Attempts to access such objects may produce an error if they execute in definer security context. (For information about security context, see Section 23.6, “Access Control for Stored Programs and Views”.)

The privilege changes take effect as indicated in Section 6.2.8, “When Privilege Changes Take Effect”.

13.7.1.8 REVOKE Syntax

```
REVOKE
    priv_type [(column_list)]
      [, priv_type [(column_list)]] ...
    ON [object_type] priv_level
    FROM user_or_role [, user_or_role] ...

REVOKE ALL [PRIVILEGES], GRANT OPTION
    FROM user_or_role [, user_or_role] ...

REVOKE PROXY ON user_or_role
    FROM user_or_role [, user_or_role] ...

REVOKE role [, role ] ...
    FROM user_or_role [, user_or_role ] ...

user_or_role: {
    user
  | role
}

user:
    (see Section 6.2.4, “Specifying Account Names”)

role:
    (see Section 6.2.5, “Specifying Role Names”.
```

The `REVOKE` statement enables system administrators to revoke privileges and roles, which can be revoked from user accounts and roles.

For information about roles, see Section 6.3.4, “Using Roles”.

When the `read_only` system variable is enabled, `REVOKE` requires the `SUPER` privilege in addition to any other required privileges described in the following discussion.

`REVOKE` either succeeds for all named users and roles or rolls back and has no effect if any error occurs. The statement is written to the binary log only if it succeeds for all named users and roles.

Each account name uses the format described in Section 6.2.4, “Specifying Account Names”. Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
REVOKE INSERT ON *.* FROM 'jeffrey'@'localhost';
REVOKE 'role1', 'role2' FROM 'user1'@'localhost', 'user2'@'localhost';
REVOKE SELECT ON world.* FROM 'role3';
```

The host name part of the account or role name, if omitted, defaults to `'%'`.

For details on the levels at which privileges exist, the permissible *`priv_type`*, *`priv_level`*, and *`object_type`* values, and the syntax for specifying users and passwords, see Section 13.7.1.6, “GRANT Syntax”

To use the first `REVOKE` syntax, you must have the `GRANT OPTION` privilege, and you must have the privileges that you are revoking.

To revoke all privileges, use the second syntax, which drops all global, database, table, column, and routine privileges for the named users or roles:

```
REVOKE ALL PRIVILEGES, GRANT OPTION
  FROM user_or_role [, user_or_role] ...
```

`REVOKE ALL PRIVILEGES, GRANT OPTION` does not revoke any roles.

To use this `REVOKE` syntax, you must have the global `CREATE USER` privilege, or the `UPDATE` privilege for the `mysql` database.

The syntax for which the `REVOKE` keyword is followed by one or more role names takes a `FROM` clause indicating one or more users or roles from which to revoke the roles.

A revoked role immediately affects any user account from which it was revoked, such that within any current session for the account, its privileges are adjusted for the next statement executed.

Revoking a role revokes the role itself, not the privileges that it represents. If an account is granted a role that includes a given privilege, and is also granted the privilege explicitly or another role that includes the privilege, the account still is granted that privilege after the first role is revoked. For example, if an account is granted two roles that each include `SELECT`, the account still can select after either role is revoked.

`REVOKE ALL ON *.*` (at the global level) revokes all granted static global privileges and all granted dynamic privileges.

User accounts and roles from which privileges and roles are to be revoked must exist, but the roles to be revoked need not be currently granted to them.

`REVOKE` removes privileges, but does not drop `mysql.user` table entries. To remove a user account entirely, use `DROP USER`. See Section 13.7.1.5, “DROP USER Syntax”.

If the grant tables hold privilege rows that contain mixed-case database or table names and the `lower_case_table_names` system variable is set to a nonzero value, `REVOKE` cannot be used to revoke these privileges. It will be necessary to manipulate the grant tables directly. (`GRANT` will not create such rows when `lower_case_table_names` is set, but such rows might have been created prior to setting the variable.)

When successfully executed from the `mysql` program, `REVOKE` responds with `Query OK, 0 rows affected`. To determine what privileges remain after the operation, use `SHOW GRANTS`. See Section 13.7.5.21, “SHOW GRANTS Syntax”.

13.7.1.9 SET DEFAULT ROLE Syntax

```
SET DEFAULT ROLE
    {NONE | ALL | role [, role ] ...}
    TO user [, user ] ...
```

For each *`user`* named immediately after the `TO` keyword, this statement defines which roles become active when the user connects to the server and authenticates, or when the user executes the `SET ROLE DEFAULT` statement during a session.

`SET DEFAULT ROLE` is alternative syntax for `ALTER USER ... DEFAULT ROLE` (see Section 13.7.1.1, “ALTER USER Syntax”). However, `ALTER USER` can set the default for only a single user, whereas `SET DEFAULT ROLE` can set the default for multiple users.

`SET DEFAULT ROLE` requires these privileges:

- To set the default roles for another user: The global `CREATE USER` privilege, or the `UPDATE` privilege for the `mysql.default_roles` system table.
- To set the default roles for yourself: No special privileges, as long as the roles you want as the default have been granted to you.

Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
SET DEFAULT ROLE administrator, developer TO 'joe'@'10.0.0.1';
```

The host name part of the role name, if omitted, defaults to `'%'`.

The clause following the `DEFAULT ROLE` keywords permits these values:

- `NONE`: Set the default to `NONE` (no roles).
- `ALL`: Set the default to all roles granted to the account.
- *`role [, role ] ...`*: Set the default to the named roles, which need not exist or be granted to the account at the time `SET DEFAULT ROLE` is executed.

Note

`SET DEFAULT ROLE` and `SET ROLE DEFAULT` are different statements. The first defines which roles to activate by default within sessions for an account. The second sets the active roles within the current session to the current account default roles.

13.7.1.10 SET PASSWORD Syntax

```
SET PASSWORD [FOR user] = password_option

password_option: {
    PASSWORD('auth_string')
  | 'auth_string'
}
```

The `SET PASSWORD` statement assigns a password to a MySQL user account, specified as either a cleartext (unencrypted) or encrypted value:

- *`'auth_string'`* represents a cleartext password.
- *`'hash_string'`* represents an encrypted password.

Note

- `SET PASSWORD ... = PASSWORD('auth_string')` syntax is deprecated as of MySQL 5.7.6 and will be removed in a future MySQL release.
- `SET PASSWORD ... = 'auth_string'` syntax is not deprecated, but `ALTER USER` is now the preferred statement for assigning passwords. For example:

```
ALTER USER user IDENTIFIED BY 'auth_string';
```

`SET PASSWORD` can be used with or without an explicitly named user account:

- With a `FOR user` clause, the statement sets the password for the named account, which must exist:

  ```
  SET PASSWORD FOR 'jeffrey'@'localhost' = password_option;
  ```

  In this case, you must have the `UPDATE` privilege for the `mysql` database.

- With no `FOR user` clause, the statement sets the password for the current user:

  ```
  SET PASSWORD = password_option;
  ```

  Any client who connects to the server using a nonanonymous account can change the password for that account. To see which account the server authenticated you as, invoke the `CURRENT_USER()` function:

  ```
  SELECT CURRENT_USER();
  ```

When the `read_only` system variable is enabled, `SET PASSWORD` requires the `SUPER` privilege in addition to any other required privileges.

If a `FOR user` clause is given, the account name uses the format described in Section 6.2.4, “Specifying Account Names”. The `user` value should be given as `'user_name'@'host_name'`, where `'user_name'` and `'host_name'` are exactly as listed in the `User` and `Host` columns of the account's `mysql.user` table row. The host name part of the account name, if omitted, defaults to `'%'`. For example, to set the password for an account with `User` and `Host` column values of `'bob'` and `'%.example.org'`, write the statement like this:

```
SET PASSWORD FOR 'bob'@'%.example.org' = PASSWORD('auth_string');
```

The password can be specified in these ways:

- Using the `PASSWORD()` function (deprecated as of MySQL 5.7.6)

  The `'auth_string'` function argument is the cleartext (unencrypted) password. `PASSWORD()` hashes the password and returns the encrypted password string for storage in the `mysql.user` account row.

  The `PASSWORD()` function hashes the password using the hashing method determined by the value of the `old_passwords` system variable value. If `SET PASSWORD` rejects the hashed password value returned by `PASSWORD()` as not being in the correct format, it may be necessary to change `old_passwords` to change the hashing method. For example, if the account uses the `mysql_native_password` plugin, the `old_passwords` value must be 0:

  ```
  SET old_passwords = 0;
  SET PASSWORD FOR 'jeffrey'@'localhost' = PASSWORD('mypass');
  ```

  If the `old_passwords` value differs from that required by the authentication plugin, the hashed password value returned by `PASSWORD()` is not acceptable for that plugin, and attempts to set the password produce an error. Permitted `old_passwords` values are described later in this section.

- Using a string without `PASSWORD()`

  `SET PASSWORD` interprets the string as a cleartext string and hashes it appropriately for the account authentication plugin before storing it in the `mysql.user` account row.

For more information about setting passwords, see Section 6.3.7, “Assigning Account Passwords”.

The following table shows the permitted values of `old_passwords`, the password hashing method for each value, and which authentication plugins use passwords hashed with each method.

Value	Password Hashing Method	Associated Authentication Plugin
0	MySQL 4.1 native hashing	`mysql_native_password`
2	SHA-256 hashing	`sha256_password`

Important

Under some circumstances, `SET PASSWORD` may be recorded in server logs or on the client side in a history file such as `~/.mysql_history`, which means that cleartext passwords may be read by anyone having read access to that information. For information about the conditions under which this occurs for the server logs and how to control it, see Section 6.1.2.3, “Passwords and Logging”. For similar information about client-side logging, see Section 4.5.1.3, “mysql Logging”.

If you are using MySQL Replication, be aware that, currently, a password used by a replication slave as part of a `CHANGE MASTER TO` statement is effectively limited to 32 characters in length; if the password is longer, any excess characters are truncated. This is not due to any limit imposed by the MySQL Server generally, but rather is an issue specific to MySQL Replication. (For more information, see Bug #43439.)

13.7.1.11 SET ROLE Syntax

```
SET ROLE
    {
        DEFAULT
      | NONE
      | ALL
      | ALL EXCEPT role [, role ] ...
      | role [, role ] ...
    }
```

`SET ROLE` modifies the current user's effective privileges within the current session by specifying which of its granted roles are active.

Any privileges that the user has been granted directly (rather than through roles) remain unaffected by changes to the active roles.

Each role name uses the format described in Section 6.2.5, “Specifying Role Names”. For example:

```
SET ROLE DEFAULT;
SET ROLE 'role1', 'role2';
SET ROLE ALL;
SET ROLE ALL EXCEPT 'role1', 'role2';
```

The host name part of the role name, if omitted, defaults to `'%'`.

The statement permits these role specifiers:

- `DEFAULT`: Set active the account default roles.

  If a user executes `SET ROLE DEFAULT` during a session, an error occurs if any default role cannot be activated (for example, if it does not exist or is not granted to the user). In this case, the current active roles are not changed.

  The server executes `SET ROLE DEFAULT` implicitly when a user connects to the server and authenticates successfully, to initialize the session active roles. In this case, the current active roles

are those default roles that can be activated. The server writes warnings to its error log for default roles that cannot be activated, but the client receives no warnings.

- `NONE`: Set the active roles to `NONE` (no active roles).
- `ALL`: Set active all roles granted to the account.
- `ALL EXCEPT role_list`: Set active all roles granted to the account except those named. The named roles need not exist or be granted to the account.
- `role [, role ] ...`: Set active the named roles, which must be granted to the account.

Note

`SET DEFAULT ROLE` and `SET ROLE DEFAULT` are different statements. The first defines which roles to activate by default within sessions for an account. The second sets the active roles within the current session to the current account default roles.

13.7.2 Table Maintenance Statements

13.7.2.1 ANALYZE TABLE Syntax

```
ANALYZE [NO_WRITE_TO_BINLOG | LOCAL]
    TABLE tbl_name [, tbl_name] ...

ANALYZE [NO_WRITE_TO_BINLOG | LOCAL]
    TABLE tbl_name
    UPDATE HISTOGRAM ON col_name [, col_name] ...
        WITH N BUCKETS

ANALYZE [NO_WRITE_TO_BINLOG | LOCAL]
    TABLE tbl_name
    DROP HISTOGRAM ON col_name [, col_name] ...
```

`ANALYZE TABLE` generates table statistics:

- `ANALYZE TABLE` without either `HISTOGRAM` clause performs a key distribution analysis and stores the distribution for the named table or tables. For `MyISAM` tables, `ANALYZE TABLE` for key distribution analysis is equivalent to using `myisamchk --analyze`.
- `ANALYZE TABLE` with the `UPDATE HISTOGRAM` clause generates histogram statistics for the named table columns and stores them in the data dictionary. Only one table name is permitted for this syntax.
- `ANALYZE TABLE` with the `DROP HISTOGRAM` clause removes histogram statistics for the named table columns from the data dictionary. Only one table name is permitted for this syntax.

The `UPDATE HISTOGRAM` and `DROP HISTOGRAM` syntaxes were added in MySQL 8.0.2.

Note

If the `innodb_read_only` system variable is enabled, `ANALYZE TABLE` may fail because it cannot update statistics tables in the data dictionary, which use `InnoDB`. For `ANALYZE TABLE` operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a `MyISAM` table). To obtain the updated distribution statistics, set `information_schema_stats=LATEST`.

This statement requires `SELECT` and `INSERT` privileges for the table.

`ANALYZE TABLE` works with `InnoDB`, `NDB`, and `MyISAM` tables. It does not work with views.

`ANALYZE TABLE` is supported for partitioned tables, and you can use `ALTER TABLE ... ANALYZE PARTITION` to analyze one or more partitions; for more information, see Section 13.1.7, “ALTER TABLE Syntax”, and Section 22.3.4, “Maintenance of Partitions”.

During the analysis, the table is locked with a read lock for `InnoDB` and `MyISAM`.

By default, the server writes `ANALYZE TABLE` statements to the binary log so that they replicate to replication slaves. To suppress logging, specify the optional `NO_WRITE_TO_BINLOG` keyword or its alias `LOCAL`.

- ANALYZE TABLE Output
- Key Distribution Analysis
- Histogram Statistics Analysis

ANALYZE TABLE Output

`ANALYZE TABLE` returns a result set with the columns shown in the following table.

Column	Value
`Table`	The table name
`Op`	`analyze` or `histogram`
`Msg_type`	`status`, `error`, `info`, `note`, or `warning`
`Msg_text`	An informational message

Key Distribution Analysis

`ANALYZE TABLE` without either `HISTOGRAM` clause performs a key distribution analysis and stores the distribution for the table or tables. Any existing histogram statistics remain unaffected.

If the table has not changed since the last key distribution analysis, the table is not analyzed again.

MySQL uses the stored key distribution to decide the order in which tables should be joined for joins on something other than a constant. In addition, key distributions can be used when deciding which indexes to use for a specific table within a query.

For more information on how key distribution analysis works within `InnoDB`, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters” and Section 15.6.11.3, “Estimating ANALYZE TABLE Complexity for InnoDB Tables”. Also see Section 15.8.1.7, “Limits on InnoDB Tables”. In particular, when you enable the `innodb_stats_persistent` option, you must run `ANALYZE TABLE` after loading substantial data into an `InnoDB` table, or creating a new index for one.

To check the stored key distribution cardinality, use the `SHOW INDEX` statement or the `INFORMATION_SCHEMA.STATISTICS` table. See Section 13.7.5.22, “SHOW INDEX Syntax”, and Section 24.21, “The INFORMATION_SCHEMA STATISTICS Table”.

Histogram Statistics Analysis

`ANALYZE TABLE` with the `HISTOGRAM` clauses enables management of histogram statistics for table column values. For information about histogram statistics, see Section 8.9.6, “Optimizer Statistics”.

These histogram operations are available:

- `ANALYZE TABLE` with an `UPDATE HISTOGRAM` clause generates histogram statistics for the named table columns and stores them in the data dictionary. Only one table name is permitted for this syntax.

The `WITH N BUCKETS` clauses specifies the number of buckets for the histogram. The value of *`N`* must be an integer in the range from 1 to 1024.

- `ANALYZE TABLE` with a `DROP HISTOGRAM` clause removes histogram statistics for the named table columns from the data dictionary. Only one table name is permitted for this syntax.

Stored histogram management statements affect only the named columns. Consider these statements:

```
ANALYZE TABLE t UPDATE HISTOGRAM ON c1, c2, c3 WITH 10 BUCKETS;
ANALYZE TABLE t UPDATE HISTOGRAM ON c1, c3 WITH 10 BUCKETS;
ANALYZE TABLE t DROP HISTOGRAM ON c2;
```

The first statement updates the histograms for columns `c1`, `c2`, and `c3`, replacing any existing histograms for those columns. The second statement updates the histograms for `c1` and `c3`, leaving the `c2` histogram unaffected. The third statement removes the histogram for `c2`, leaving those for `c1` and `c3` unaffected.

Histogram generation is not supported for encrypted tables (to avoid exposing data in the statistics) or `TEMPORARY` tables.

Histogram generation applies to columns of all data types except geometry types (spatial data) and `JSON`.

Histograms can be generated for stored and virtual generated columns.

Histograms cannot be generated for columns that are covered by single-column unique indexes.

Histogram management statements attempt to perform as much of the requested operation as possible, and report diagnostic messages for the remainder. For example, if an `UPDATE HISTOGRAM` statement names multiple columns, but some of them do not exist or have an unsupported data type, histograms are generated for the other columns, and messages are produced for the invalid columns.

The `histogram_generation_max_mem_size` system variable controls the maximum amount of memory available for histogram generation. The global and session values may be set at runtime. The `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege is required, even for the session value.

For information about memory allocations performed for histogram generation, monitor the Performance Schema `memory/sql/histograms` instrument. See Section 25.11.15.10, “Memory Summary Tables”.

Histograms are affected by these DDL statements:

- `DROP TABLE` removes histograms for columns in the dropped table.
- `DROP DATABASE` removes histograms for any table in the dropped database because the statement drops all tables in the database.
- `RENAME TABLE` does not remove histograms. Instead, it renames histograms for the renamed table to be associated with the new table name.
- `ALTER TABLE` statements that remove or modify a column remove histograms for that column.
- `ALTER TABLE ... CONVERT TO CHARACTER SET` removes histograms for character columns because they are affected by the change of character set. Histograms for noncharacter columns remain unaffected.

13.7.2.2 CHECK TABLE Syntax

```
CHECK TABLE tbl_name [, tbl_name] ... [option] ...
```

```
option = {
    FOR UPGRADE
  | QUICK
  | FAST
  | MEDIUM
  | EXTENDED
  | CHANGED
}
```

`CHECK TABLE` checks a table or tables for errors. For `MyISAM` tables, the key statistics are updated as well. `CHECK TABLE` can also check views for problems, such as tables that are referenced in the view definition that no longer exist.

To check a table, you must have some privilege for it.

`CHECK TABLE` works for `InnoDB`, `MyISAM`, `ARCHIVE`, and `CSV` tables.

Before running `CHECK TABLE` on `InnoDB` tables, see CHECK TABLE Usage Notes for InnoDB Tables.

`CHECK TABLE` is supported for partitioned tables, and you can use `ALTER TABLE ... CHECK PARTITION` to check one or more partitions; for more information, see Section 13.1.7, “ALTER TABLE Syntax”, and Section 22.3.4, “Maintenance of Partitions”.

`CHECK TABLE` ignores virtual generated columns that are not indexed.

- CHECK TABLE Output
- Checking Version Compatibility
- Checking Data Consistency
- CHECK TABLE Usage Notes for InnoDB Tables

CHECK TABLE Output

`CHECK TABLE` returns a result set with the columns shown in the following table.

Column	Value
`Table`	The table name
`Op`	Always `check`
`Msg_type`	`status`, `error`, `info`, `note`, or `warning`
`Msg_text`	An informational message

The statement might produce many rows of information for each checked table. The last row has a `Msg_type` value of `status` and the `Msg_text` normally should be `OK`. For a `MyISAM` table, if you don't get `OK` or `Table is already up to date`, you should normally run a repair of the table. See Section 7.6, “MyISAM Table Maintenance and Crash Recovery”. `Table is already up to date` means that the storage engine for the table indicated that there was no need to check the table.

Checking Version Compatibility

The `FOR UPGRADE` option checks whether the named tables are compatible with the current version of MySQL. With `FOR UPGRADE`, the server checks each table to determine whether there have been any incompatible changes in any of the table's data types or indexes since the table was created. If not, the check succeeds. Otherwise, if there is a possible incompatibility, the server runs a full check on the table (which might take some time).

Incompatibilities might occur because the storage format for a data type has changed or because its sort order has changed. Our aim is to avoid these changes, but occasionally they are necessary to correct problems that would be worse than an incompatibility between releases.

`FOR UPGRADE` discovers these incompatibilities:

- The indexing order for end-space in `TEXT` columns for `InnoDB` and `MyISAM` tables changed between MySQL 4.1 and 5.0.
- The storage method of the new `DECIMAL` data type changed between MySQL 5.0.3 and 5.0.5.
- Changes are sometimes made to character sets or collations that require table indexes to be rebuilt. For details about such changes, see Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0”. For information about rebuilding tables, see Section 2.10.3, “Rebuilding or Repairing Tables or Indexes”.
- MySQL 8.0 does not support the `YEAR(2)` data type permitted in older versions of MySQL. For tables containing `YEAR(2)` columns, `CHECK TABLE` recommends `REPAIR TABLE`, which converts `YEAR(2)` to `YEAR(4)`.
- Trigger creation time is maintained.
- A table is reported as needing a rebuild if it contains old temporal columns in pre-5.6.4 format (`TIME`, `DATETIME`, and `TIMESTAMP` columns without support for fractional seconds precision) and the `avoid_temporal_upgrade` system variable is disabled. This helps `mysql_upgrade` detect and upgrade tables containing old temporal columns. If `avoid_temporal_upgrade` is enabled, `FOR UPGRADE` ignores the old temporal columns present in the table; consequently, `mysql_upgrade` does not upgrade them.

  To check for tables that contain such temporal columns and need a rebuild, disable `avoid_temporal_upgrade` before executing `CHECK TABLE ... FOR UPGRADE`.

Checking Data Consistency

The following table shows the other check options that can be given. These options are passed to the storage engine, which may use or ignore them.

Type	Meaning
`QUICK`	Do not scan the rows to check for incorrect links. Applies to `InnoDB` and `MyISAM` tables and views.
`FAST`	Check only tables that have not been closed properly. Applies only to `MyISAM` tables and views; ignored for `InnoDB`.
`CHANGED`	Check only tables that have been changed since the last check or that have not been closed properly. Applies only to `MyISAM` tables and views; ignored for `InnoDB`.
`MEDIUM`	Scan rows to verify that deleted links are valid. This also calculates a key checksum for the rows and verifies this with a calculated checksum for the keys. Applies only to `MyISAM` tables and views; ignored for `InnoDB`.
`EXTENDED`	Do a full key lookup for all keys for each row. This ensures that the table is 100% consistent, but takes a long time. Applies only to `MyISAM` tables and views; ignored for `InnoDB`.

If none of the options `QUICK`, `MEDIUM`, or `EXTENDED` are specified, the default check type for dynamic-format `MyISAM` tables is `MEDIUM`. This has the same result as running `myisamchk --medium-check` *`tbl_name`* on the table. The default check type also is `MEDIUM` for static-format `MyISAM` tables, unless `CHANGED` or `FAST` is specified. In that case, the default is `QUICK`. The row scan is skipped for `CHANGED` and `FAST` because the rows are very seldom corrupted.

You can combine check options, as in the following example that does a quick check on the table to determine whether it was closed properly:

```
CHECK TABLE test_table FAST QUICK;
```

Note

If `CHECK TABLE` finds no problems with a table that is marked as “corrupted” or “not closed properly”, `CHECK TABLE` may remove the mark.

If a table is corrupted, the problem is most likely in the indexes and not in the data part. All of the preceding check types check the indexes thoroughly and should thus find most errors.

To check a table that you assume is okay, use no check options or the `QUICK` option. The latter should be used when you are in a hurry and can take the very small risk that `QUICK` does not find an error in the data file. (In most cases, under normal usage, MySQL should find any error in the data file. If this happens, the table is marked as “corrupted” and cannot be used until it is repaired.)

`FAST` and `CHANGED` are mostly intended to be used from a script (for example, to be executed from `cron`) to check tables periodically. In most cases, `FAST` is to be preferred over `CHANGED`. (The only case when it is not preferred is when you suspect that you have found a bug in the `MyISAM` code.)

`EXTENDED` is to be used only after you have run a normal check but still get errors from a table when MySQL tries to update a row or find a row by key. This is very unlikely if a normal check has succeeded.

Use of `CHECK TABLE ... EXTENDED` might influence execution plans generated by the query optimizer.

Some problems reported by `CHECK TABLE` cannot be corrected automatically:

- `Found row where the auto_increment column has the value 0`.

  This means that you have a row in the table where the `AUTO_INCREMENT` index column contains the value 0. (It is possible to create a row where the `AUTO_INCREMENT` column is 0 by explicitly setting the column to 0 with an `UPDATE` statement.)

  This is not an error in itself, but could cause trouble if you decide to dump the table and restore it or do an `ALTER TABLE` on the table. In this case, the `AUTO_INCREMENT` column changes value according to the rules of `AUTO_INCREMENT` columns, which could cause problems such as a duplicate-key error.

  To get rid of the warning, execute an `UPDATE` statement to set the column to some value other than 0.

CHECK TABLE Usage Notes for InnoDB Tables

The following notes apply to `InnoDB` tables:

- If `CHECK TABLE` encounters a corrupt page, the server exits to prevent error propagation (Bug #10132). If the corruption occurs in a secondary index but table data is readable, running `CHECK TABLE` can still cause a server exit.
- If `CHECK TABLE` encounters a corrupted `DB_TRX_ID` or `DB_ROLL_PTR` field in a clustered index, `CHECK TABLE` can cause `InnoDB` to access an invalid undo log record, resulting in an MVCC-related server exit.
- If `CHECK TABLE` encounters errors in `InnoDB` tables or indexes, it reports an error, and usually marks the index and sometimes marks the table as corrupted, preventing further use of the index or table. Such errors include an incorrect number of entries in a secondary index or incorrect links.
- If `CHECK TABLE` finds an incorrect number of entries in a secondary index, it reports an error but does not cause a server exit or prevent access to the file.

- `CHECK TABLE` surveys the index page structure, then surveys each key entry. It does not validate the key pointer to a clustered record or follow the path for `BLOB` pointers.
- When an `InnoDB` table is stored in its own `.ibd` file, the first 3 pages of the `.ibd` file contain header information rather than table or index data. The `CHECK TABLE` statement does not detect inconsistencies that affect only the header data. To verify the entire contents of an `InnoDB .ibd` file, use the `innochecksum` command.
- When running `CHECK TABLE` on large `InnoDB` tables, other threads may be blocked during `CHECK TABLE` execution. To avoid timeouts, the semaphore wait threshold (600 seconds) is extended by 2 hours (7200 seconds) for `CHECK TABLE` operations. If `InnoDB` detects semaphore waits of 240 seconds or more, it starts printing `InnoDB` monitor output to the error log. If a lock request extends beyond the semaphore wait threshold, `InnoDB` aborts the process. To avoid the possibility of a semaphore wait timeout entirely, you can run `CHECK TABLE QUICK` instead of `CHECK TABLE`.
- `CHECK TABLE` functionality for `InnoDB SPATIAL` indexes includes an R-tree validity check and a check to ensure that the R-tree row count matches the clustered index.
- `CHECK TABLE` supports secondary indexes on virtual generated columns. `InnoDB` has support for secondary indexes on virtual generated columns.

13.7.2.3 CHECKSUM TABLE Syntax

```
CHECKSUM TABLE tbl_name [, tbl_name] ... [QUICK | EXTENDED]
```

`CHECKSUM TABLE` reports a checksum for the contents of a table. You can use this statement to verify that the contents are the same before and after a backup, rollback, or other operation that is intended to put the data back to a known state.

This statement requires the `SELECT` privilege for the table.

This statement is not supported for views. If you run `CHECKSUM TABLE` against a view, the `Checksum` value is always `NULL`, and a warning is returned.

For a nonexistent table, `CHECKSUM TABLE` returns `NULL` and generates a warning.

During the checksum operation, the table is locked with a read lock for `InnoDB` and `MyISAM`.

Performance Considerations

By default, the entire table is read row by row and the checksum is calculated. For large tables, this could take a long time, thus you would only perform this operation occasionally. This row-by-row calculation is what you get with the `EXTENDED` clause, with `InnoDB` and all other storage engines other than `MyISAM`, and with `MyISAM` tables not created with the `CHECKSUM=1` clause.

For `MyISAM` tables created with the `CHECKSUM=1` clause, `CHECKSUM TABLE` or `CHECKSUM TABLE ... QUICK` returns the “live” table checksum that can be returned very fast. If the table does not meet all these conditions, the `QUICK` method returns `NULL`. See Section 13.1.15, “CREATE TABLE Syntax” for the syntax of the `CHECKSUM` clause.

The checksum value depends on the table row format. If the row format changes, the checksum also changes. For example, the storage format for temporal types such as `TIME`, `DATETIME`, and `TIMESTAMP` changed in MySQL 5.6 prior to MySQL 5.6.5, so if a 5.5 table is upgraded to MySQL 5.6, the checksum value may change.

Important

If the checksums for two tables are different, then it is almost certain that the tables are different in some way. However, because the hashing function used

by `CHECKSUM TABLE` is not guaranteed to be collision-free, there is a slight chance that two tables which are not identical can produce the same checksum.

13.7.2.4 OPTIMIZE TABLE Syntax

```
OPTIMIZE [NO_WRITE_TO_BINLOG | LOCAL]
    TABLE tbl_name [, tbl_name] ...
```

`OPTIMIZE TABLE` reorganizes the physical storage of table data and associated index data, to reduce storage space and improve I/O efficiency when accessing the table. The exact changes made to each table depend on the storage engine used by that table.

Use `OPTIMIZE TABLE` in these cases, depending on the type of table:

- After doing substantial insert, update, or delete operations on an `InnoDB` table that has its own .ibd file because it was created with the `innodb_file_per_table` option enabled. The table and indexes are reorganized, and disk space can be reclaimed for use by the operating system.
- After doing substantial insert, update, or delete operations on columns that are part of a `FULLTEXT` index in an `InnoDB` table. Set the configuration option `innodb_optimize_fulltext_only=1` first. To keep the index maintenance period to a reasonable time, set the `innodb_ft_num_word_optimize` option to specify how many words to update in the search index, and run a sequence of `OPTIMIZE TABLE` statements until the search index is fully updated.
- After deleting a large part of a `MyISAM` or `ARCHIVE` table, or making many changes to a `MyISAM` or `ARCHIVE` table with variable-length rows (tables that have `VARCHAR`, `VARBINARY`, `BLOB`, or `TEXT` columns). Deleted rows are maintained in a linked list and subsequent `INSERT` operations reuse old row positions. You can use `OPTIMIZE TABLE` to reclaim the unused space and to defragment the data file. After extensive changes to a table, this statement may also improve performance of statements that use the table, sometimes significantly.

This statement requires `SELECT` and `INSERT` privileges for the table.

`OPTIMIZE TABLE` works for `InnoDB`, `MyISAM`, and `ARCHIVE` tables.

By default, `OPTIMIZE TABLE` does *not* work for tables created using any other storage engine and returns a result indicating this lack of support. You can make `OPTIMIZE TABLE` work for other storage engines by starting `mysqld` with the `--skip-new` option. In this case, `OPTIMIZE TABLE` is just mapped to `ALTER TABLE`.

This statement does not work with views.

`OPTIMIZE TABLE` is supported for partitioned tables. For information about using this statement with partitioned tables and table partitions, see Section 22.3.4, “Maintenance of Partitions”.

By default, the server writes `OPTIMIZE TABLE` statements to the binary log so that they replicate to replication slaves. To suppress logging, specify the optional `NO_WRITE_TO_BINLOG` keyword or its alias `LOCAL`.

- OPTIMIZE TABLE Output
- InnoDB Details
- MyISAM Details
- Other Considerations

OPTIMIZE TABLE Output

`OPTIMIZE TABLE` returns a result set with the columns shown in the following table.

Column	Value
`Table`	The table name
`Op`	Always `optimize`
`Msg_type`	`status`, `error`, `info`, `note`, or `warning`
`Msg_text`	An informational message

`OPTIMIZE TABLE` table catches and throws any errors that occur while copying table statistics from the old file to the newly created file. For example. if the user ID of the owner of the `.MYD` or `.MYI` file is different from the user ID of the `mysqld` process, `OPTIMIZE TABLE` generates a "cannot change ownership of the file" error unless `mysqld` is started by the `root` user.

InnoDB Details

For `InnoDB` tables, `OPTIMIZE TABLE` is mapped to `ALTER TABLE ... FORCE`, which rebuilds the table to update index statistics and free unused space in the clustered index. This is displayed in the output of `OPTIMIZE TABLE` when you run it on an `InnoDB` table, as shown here:

```
mysql> OPTIMIZE TABLE foo;
+----------+----------+----------+-------------------------------------------------------------------+
| Table    | Op       | Msg_type | Msg_text                                                          |
+----------+----------+----------+-------------------------------------------------------------------+
| test.foo | optimize | note     | Table does not support optimize, doing recreate + analyze instead |
| test.foo | optimize | status   | OK                                                                |
+----------+----------+----------+-------------------------------------------------------------------+
```

`OPTIMIZE TABLE` uses online DDL for regular and partitioned `InnoDB` tables, which reduces downtime for concurrent DML operations. The table rebuild triggered by `OPTIMIZE TABLE` and performed under the cover by `ALTER TABLE ... FORCE` is completed in place. An exclusive table lock is only taken briefly during the prepare phase and the commit phase of the operation. During the prepare phase, metadata is updated and an intermediate table is created. During the commit phase, table metadata changes are committed.

`OPTIMIZE TABLE` rebuilds the table using the table copy method under the following conditions:

- When the `old_alter_table` system variable is enabled.
- When the `mysqld --skip-new` option is enabled.

`OPTIMIZE TABLE` using online DDL is not supported for `InnoDB` tables that contain `FULLTEXT` indexes. The table copy method is used instead.

`InnoDB` stores data using a page-allocation method and does not suffer from fragmentation in the same way that legacy storage engines (such as `MyISAM`) will. When considering whether or not to run optimize, consider the workload of transactions that your server will process:

- Some level of fragmentation is expected. `InnoDB` only fills pages 93% full, to leave room for updates without having to split pages.
- Delete operations might leave gaps that leave pages less filled than desired, which could make it worthwhile to optimize the table.
- Updates to rows usually rewrite the data within the same page, depending on the data type and row format, when sufficient space is available. See Section 15.9.1.5, “How Compression Works for InnoDB Tables” and Section 15.10.1, “Overview of InnoDB Row Storage”.
- High-concurrency workloads might leave gaps in indexes over time, as `InnoDB` retains multiple versions of the same data due through its MVCC mechanism. See Section 15.3, “InnoDB Multi-Versioning”.

MyISAM Details

For MyISAM tables, OPTIMIZE TABLE works as follows:

1. If the table has deleted or split rows, repair the table.
2. If the index pages are not sorted, sort them.
3. If the table's statistics are not up to date (and the repair could not be accomplished by sorting the index), update them.

Other Considerations

OPTIMIZE TABLE is performed online for regular and partitioned InnoDB tables. Otherwise, MySQL locks the table during the time OPTIMIZE TABLE is running.

OPTIMIZE TABLE does not sort R-tree indexes, such as spatial indexes on POINT columns. (Bug #23578)

13.7.2.5 REPAIR TABLE Syntax

```
REPAIR [NO_WRITE_TO_BINLOG | LOCAL]
    TABLE tbl_name [, tbl_name] ...
    [QUICK] [EXTENDED] [USE_FRM]
```

REPAIR TABLE repairs a possibly corrupted table, for certain storage engines only. For MyISAM tables, it has the same effect as myisamchk --recover *tbl_name* by default.

This statement requires SELECT and INSERT privileges for the table.

REPAIR TABLE works for MyISAM, ARCHIVE, and CSV tables. See Section 16.2, “The MyISAM Storage Engine” Section 16.5, “The ARCHIVE Storage Engine”, and Section 16.4, “The CSV Storage Engine”. This statement does not work with views.

You can use ALTER TABLE ... REPAIR PARTITION to repair one or more partitions; for more information, see Section 13.1.7, “ALTER TABLE Syntax”, and Section 22.3.4, “Maintenance of Partitions”.

Although normally you should never have to run REPAIR TABLE, if disaster strikes, this statement is very likely to get back all your data from a MyISAM table. If your tables become corrupted often, try to find the reason for it, to eliminate the need to use REPAIR TABLE. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”, and Section 16.2.4, “MyISAM Table Problems”.

Caution

Make a backup of a table before performing a table repair operation; under some circumstances the operation might cause data loss. Possible causes include but are not limited to file system errors. See Chapter 7, *Backup and Recovery*.

Warning

If the server crashes during a REPAIR TABLE operation, it is essential after restarting it that you immediately execute another REPAIR TABLE statement for the table before performing any other operations on it. In the worst case, you might have a new clean index file without information about the data file, and then the next operation you perform could overwrite the data file. This is an unlikely but possible scenario that underscores the value of making a backup first.

By default, the server writes `REPAIR TABLE` statements to the binary log so that they replicate to replication slaves. To suppress logging, specify the optional `NO_WRITE_TO_BINLOG` keyword or its alias `LOCAL`.

Important

In the event that a table on the master becomes corrupted and you run `REPAIR TABLE` on it, any resulting changes to the original table are *not* propagated to slaves.

- REPAIR TABLE Output
- Table Repair Considerations

REPAIR TABLE Output

`REPAIR TABLE` returns a result set with the columns shown in the following table.

Column	Value
`Table`	The table name
`Op`	Always `repair`
`Msg_type`	`status`, `error`, `info`, `note`, or `warning`
`Msg_text`	An informational message

The `REPAIR TABLE` statement might produce many rows of information for each repaired table. The last row has a `Msg_type` value of `status` and `Msg_test` normally should be `OK`. For a `MyISAM` table, if you do not get `OK`, you should try repairing it with `myisamchk --safe-recover`. (`REPAIR TABLE` does not implement all the options of `myisamchk`. With `myisamchk --safe-recover`, you can also use options that `REPAIR TABLE` does not support, such as `--max-record-length`.)

`REPAIR TABLE` table catches and throws any errors that occur while copying table statistics from the old corrupted file to the newly created file. For example. if the user ID of the owner of the `.MYD` or `.MYI` file is different from the user ID of the `mysqld` process, `REPAIR TABLE` generates a "cannot change ownership of the file" error unless `mysqld` is started by the `root` user.

Table Repair Considerations

If you use the `QUICK` option, `REPAIR TABLE` tries to repair only the index file, and not the data file. This type of repair is like that done by `myisamchk --recover --quick`.

If you use the `EXTENDED` option, MySQL creates the index row by row instead of creating one index at a time with sorting. This type of repair is like that done by `myisamchk --safe-recover`.

The `USE_FRM` option is available for use if the `.MYI` index file is missing or if its header is corrupted. This option tells MySQL not to trust the information in the `.MYI` file header and to re-create it using information from the data dictionary. This kind of repair cannot be done with `myisamchk`.

Caution

Use the `USE_FRM` option *only* if you cannot use regular `REPAIR` modes. Telling the server to ignore the `.MYI` file makes important table metadata stored in the `.MYI` unavailable to the repair process, which can have deleterious consequences:

- The current `AUTO_INCREMENT` value is lost.
- The link to deleted records in the table is lost, which means that free space for deleted records will remain unoccupied thereafter.

- The `.MYI` header indicates whether the table is compressed. If the server ignores this information, it cannot tell that a table is compressed and repair can cause change or loss of table contents. This means that `USE_FRM` should not be used with compressed tables. That should not be necessary, anyway: Compressed tables are read only, so they should not become corrupt.

If `USE_FRM` is *not* used, `REPAIR TABLE` checks the table to see whether an upgrade is required. If so, it performs the upgrade, following the same rules as `CHECK TABLE ... FOR UPGRADE`. See Section 13.7.2.2, “CHECK TABLE Syntax”, for more information.

`REPAIR TABLE` checks the table to see whether an upgrade is required. If so, it performs the upgrade, following the same rules as `CHECK TABLE ... FOR UPGRADE`. See Section 13.7.2.2, “CHECK TABLE Syntax”, for more information.

`REPAIR TABLE` upgrades a table if it contains old temporal columns in pre-5.6.4 format (`TIME`, `DATETIME`, and `TIMESTAMP` columns without support for fractional seconds precision) and the `avoid_temporal_upgrade` system variable is disabled. If `avoid_temporal_upgrade` is enabled, `REPAIR TABLE` ignores the old temporal columns present in the table and does not upgrade them.

To upgrade tables that contain such temporal columns, disable `avoid_temporal_upgrade` before executing `REPAIR TABLE`.

You may be able to increase `REPAIR TABLE` performance by setting certain system variables. See Section 8.6.3, “Optimizing REPAIR TABLE Statements”.

13.7.3 Component, Plugin, and User-Defined Function Statements

13.7.3.1 CREATE FUNCTION Syntax for User-Defined Functions

```
CREATE [AGGREGATE] FUNCTION function_name
    RETURNS {STRING|INTEGER|REAL|DECIMAL}
    SONAME shared_library_name
```

A user-defined function (UDF) is a way to extend MySQL with a new function that works like a native (built-in) MySQL function such as `ABS()` or `CONCAT()`.

`function_name` is the name that should be used in SQL statements to invoke the function. The `RETURNS` clause indicates the type of the function's return value. `DECIMAL` is a legal value after `RETURNS`, but currently `DECIMAL` functions return string values and should be written like `STRING` functions.

`shared_library_name` is the base name of the shared library file that contains the code that implements the function. The file must be located in the plugin directory. This directory is given by the value of the `plugin_dir` system variable. For more information, see Section 28.4.2.5, “UDF Compiling and Installing”.

To create a function, you must have the `INSERT` privilege for the `mysql` database. This is necessary because `CREATE FUNCTION` adds a row to the `mysql.func` system table that records the function's name, type, and shared library name. If you do not have this table, you should run the `mysql_upgrade` command to create it. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

As of MySQL 8.0.2, UDFs registered using `CREATE FUNCTION` are listed in the Performance Schema `user_defined_functions` table; see Section 25.11.16.4, “The user_defined_functions Table”.

An active function is one that has been loaded with `CREATE FUNCTION` and not removed with `DROP FUNCTION`. All active functions are reloaded each time the server starts, unless you start `mysqld`

with the `--skip-grant-tables` option. In this case, UDF initialization is skipped and UDFs are unavailable.

For instructions on writing user-defined functions, see Section 28.4.2, "Adding a New User-Defined Function". For the UDF mechanism to work, functions must be written in C or C++ (or another language that can use C calling conventions), your operating system must support dynamic loading and you must have compiled `mysqld` dynamically (not statically).

An `AGGREGATE` function works exactly like a native MySQL aggregate (summary) function such as `SUM` or `COUNT()`. For `AGGREGATE` to work, your `mysql.func` table must contain a `type` column. If your `mysql.func` table does not have this column, you should run the `mysql_upgrade` program to create it (see Section 4.4.5, "`mysql_upgrade` — Check and Upgrade MySQL Tables").

Note

To upgrade the shared library associated with a UDF, issue a `DROP FUNCTION` statement, upgrade the shared library, and then issue a `CREATE FUNCTION` statement. If you upgrade the shared library first and then use `DROP FUNCTION`, the server may crash.

13.7.3.2 DROP FUNCTION Syntax

```
DROP FUNCTION function_name
```

This statement drops the user-defined function (UDF) named *`function_name`*.

To drop a function, you must have the `DELETE` privilege for the `mysql` database. This is because `DROP FUNCTION` removes a row from the `mysql.func` system table that records the function's name, type, and shared library name.

Note

To upgrade the shared library associated with a UDF, issue a `DROP FUNCTION` statement, upgrade the shared library, and then issue a `CREATE FUNCTION` statement. If you upgrade the shared library first and then use `DROP FUNCTION`, the server may crash.

`DROP FUNCTION` is also used to drop stored functions (see Section 13.1.23, "DROP PROCEDURE and DROP FUNCTION Syntax").

13.7.3.3 INSTALL COMPONENT Syntax

```
INSTALL COMPONENT component_name [, component_name ] ...
```

This statement installs one or more server components, which become active immediately. A component provides services that are available to the server and other components. (With respect to service use, the server is a component, equal to other components.) `INSTALL COMPONENT` requires the `INSERT` privilege for the `mysql.component` system table.

A loader service handles component installation and uninstallation, and also lists installed components in the `mysql.component` system table, which serves as a registry. For subsequent server restarts, any components listed in `mysql.component` are installed by the loader service during the startup sequence. This occurs even if the server is started with the `--skip-grant-tables` option.

If a component depends on services not present in the registry and you attempt to install the component without also installing the component or components that provide the services it depends on, an error occurs:

```
ERROR 3527 (HY000): Cannot satisfy dependency for service 'component1'
```

```
required by component 'component2'.
```

To avoid this, either install all components in the same statement, or install the dependent component after installing any components it depends on.

If any error occurs, the statement fails and has no effect. For example, this happens if a component name is erroneous, a named component does not exist or is already installed, or component initialization fails.

Currently, component names are URNs that begin with `file://` and for which the following file name value contains no directory names and is relative to the directory named by the `plugin_dir` system variable. Component names do not include any platform-dependent file name suffix such as `.so` or `.dll`. These details are subject to change because component name interpretation is itself performed by a service and the component infrastructure makes it possible to replace the default service implementation with alternative implementations.

Example:

```
INSTALL COMPONENT 'file://component1', 'file://component2';
```

For additional information about components, see Section 5.5, “MySQL Server Components”.

13.7.3.4 INSTALL PLUGIN Syntax

```
INSTALL PLUGIN plugin_name SONAME 'shared_library_name'
```

This statement installs a server plugin. It requires the `INSERT` privilege for the `mysql.plugin` system table.

`plugin_name` is the name of the plugin as defined in the plugin descriptor structure contained in the library file (see Section 28.2.4.2, “Plugin Data Structures”). Plugin names are not case sensitive. For maximal compatibility, plugin names should be limited to ASCII letters, digits, and underscore because they are used in C source files, shell command lines, M4 and Bourne shell scripts, and SQL environments.

`shared_library_name` is the name of the shared library that contains the plugin code. The name includes the file name extension (for example, `libmyplugin.so`, `libmyplugin.dll`, or `libmyplugin.dylib`).

The shared library must be located in the plugin directory (the directory named by the `plugin_dir` system variable). The library must be in the plugin directory itself, not in a subdirectory. By default, `plugin_dir` is the `plugin` directory under the directory named by the `pkglibdir` configuration variable, but it can be changed by setting the value of `plugin_dir` at server startup. For example, set its value in a `my.cnf` file:

```
[mysqld]
plugin_dir=/path/to/plugin/directory
```

If the value of `plugin_dir` is a relative path name, it is taken to be relative to the MySQL base directory (the value of the `basedir` system variable).

`INSTALL PLUGIN` loads and initializes the plugin code to make the plugin available for use. A plugin is initialized by executing its initialization function, which handles any setup that the plugin must perform before it can be used. When the server shuts down, it executes the deinitialization function for each plugin that is loaded so that the plugin has a chance to perform any final cleanup.

`INSTALL PLUGIN` also registers the plugin by adding a line that indicates the plugin name and library file name to the `mysql.plugin` table. At server startup, the server loads and initializes any plugin that is listed in the `mysql.plugin` table. This means that a plugin is installed with `INSTALL PLUGIN` only

once, not every time the server starts. Plugin loading at startup does not occur if the server is started with the `--skip-grant-tables` option.

A plugin library can contain multiple plugins. For each of them to be installed, use a separate `INSTALL PLUGIN` statement. Each statement names a different plugin, but all of them specify the same library name.

`INSTALL PLUGIN` causes the server to read option (`my.cnf`) files just as during server startup. This enables the plugin to pick up any relevant options from those files. It is possible to add plugin options to an option file even before loading a plugin (if the `loose` prefix is used). It is also possible to uninstall a plugin, edit `my.cnf`, and install the plugin again. Restarting the plugin this way enables it to the new option values without a server restart.

For options that control individual plugin loading at server startup, see Section 5.6.2, “Installing and Uninstalling Plugins”. If you need to load plugins for a single server startup when the `--skip-grant-tables` option is given (which tells the server not to read system tables), use the `--plugin-load` option. See Section 5.1.4, “Server Command Options”.

To remove a plugin, use the `UNINSTALL PLUGIN` statement.

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To see what plugins are installed, use the `SHOW PLUGINS` statement or query the `INFORMATION_SCHEMA.PLUGINS` table.

If you recompile a plugin library and need to reinstall it, you can use either of the following methods:

- Use `UNINSTALL PLUGIN` to uninstall all plugins in the library, install the new plugin library file in the plugin directory, and then use `INSTALL PLUGIN` to install all plugins in the library. This procedure has the advantage that it can be used without stopping the server. However, if the plugin library contains many plugins, you must issue many `INSTALL PLUGIN` and `UNINSTALL PLUGIN` statements.
- Stop the server, install the new plugin library file in the plugin directory, and restart the server.

13.7.3.5 UNINSTALL COMPONENT Syntax

```
UNINSTALL COMPONENT component_name [, component_name ] ...
```

This statement deactivates and uninstalls one or more server components. It requires the `DELETE` privilege for the `mysql.component` system table. `UNINSTALL COMPONENT` is the complement of `INSTALL COMPONENT`.

A loader service also unregisters uninstalled components from the `mysql.component` system table. As a result, they are no longer installed during the startup sequence for subsequent server restarts.

If any error occurs, the statement fails and has no effect. For example, this happens if a component name is erroneous, a named component is not installed, or cannot be uninstalled because other installed components depend on it.

For information about component naming, see Section 13.7.3.3, “INSTALL COMPONENT Syntax”.

Example:

```
UNINSTALL COMPONENT 'file://component2';
```

For additional information about components, see Section 5.5, “MySQL Server Components”.

13.7.3.6 UNINSTALL PLUGIN Syntax

```
UNINSTALL PLUGIN plugin_name
```

This statement removes an installed server plugin. It requires the `DELETE` privilege for the `mysql.plugin` system table. `UNINSTALL PLUGIN` is the complement of `INSTALL PLUGIN`.

`plugin_name` must be the name of some plugin that is listed in the `mysql.plugin` table. The server executes the plugin's deinitialization function and removes the row for the plugin from the `mysql.plugin` table, so that subsequent server restarts will not load and initialize the plugin. `UNINSTALL PLUGIN` does not remove the plugin's shared library file.

You cannot uninstall a plugin if any table that uses it is open.

Plugin removal has implications for the use of associated tables. For example, if a full-text parser plugin is associated with a `FULLTEXT` index on the table, uninstalling the plugin makes the table unusable. Any attempt to access the table results in an error. The table cannot even be opened, so you cannot drop an index for which the plugin is used. This means that uninstalling a plugin is something to do with care unless you do not care about the table contents. If you are uninstalling a plugin with no intention of reinstalling it later and you care about the table contents, you should dump the table with `mysqldump` and remove the `WITH PARSER` clause from the dumped `CREATE TABLE` statement so that you can reload the table later. If you do not care about the table, `DROP TABLE` can be used even if any plugins associated with the table are missing.

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

13.7.4 SET Syntax

The `SET` statement has several forms:

- `SET var_name = value` enables you to assign values to variables that affect the operation of the server or clients. See Section 13.7.4.1, “SET Syntax for Variable Assignment”.
- `SET CHARACTER SET` and `SET NAMES` assign values to character set and collation variables associated with the current connection to the server. See Section 13.7.4.2, “SET CHARACTER SET Syntax”, and Section 13.7.4.3, “SET NAMES Syntax”.
- `SET DEFAULT ROLE` and `SET ROLE` set the default role and current role for user accounts. See Section 13.7.1.9, “SET DEFAULT ROLE Syntax”, and Section 13.7.1.11, “SET ROLE Syntax”.
- `SET PASSWORD` assigns account passwords. See Section 13.7.1.10, “SET PASSWORD Syntax”.
- `SET TRANSACTION ISOLATION LEVEL` sets the isolation level for transaction processing. See Section 13.3.6, “SET TRANSACTION Syntax”.

13.7.4.1 SET Syntax for Variable Assignment

```
SET variable_assignment [, variable_assignment] ...

variable_assignment:
      user_var_name = expr
    | param_name = expr
    | local_var_name = expr
    | [GLOBAL | SESSION | PERSIST]
        system_var_name = expr
    | [@@global. | @@session. | @@persist. | @@]
        system_var_name = expr
```

`SET` syntax for variable assignment enables you to assign values to different types of variables that affect the operation of the server or clients:

- System variables. See Section 5.1.5, “Server System Variables”. System variables also can be set at server startup, as described in Section 5.1.6, “Using System Variables”. (To *display* system

variable names and values, use the SHOW VARIABLES statement; see Section 13.7.5.39, “SHOW VARIABLES Syntax”.)

- User-defined variables. See Section 9.4, “User-Defined Variables”.
- Stored procedure and function parameters, and stored program local variables. See Section 13.6.4, “Variables in Stored Programs”.

A SET statement that assigns variable values is not written to the binary log, so in replication scenarios it affects only the host on which you execute it. To affect all replication hosts, execute the statement on each one.

The following examples illustrate SET syntax for setting variables. They use the = assignment operator, but the := assignment operator is also permitted for this purpose.

A user variable is written as *@var_name* and is assigned an expression value as follows:

```
SET @var_name = expr;
```

Examples:

```
SET @name = 43;
SET @total_tax = (SELECT SUM(tax) FROM taxable_transactions);
```

As demonstrated by those statements, *expr* can range from simple (a literal value) to more complex (the value returned by a scalar subquery).

SET applies to parameters and local variables in the context of the stored object within which they are defined. The following procedure uses the counter local variable as a loop counter:

```
CREATE PROCEDURE p()
BEGIN
  DECLARE counter INT DEFAULT 0;
  WHILE counter < 10 DO
    -- ... do work ...
    SET counter = counter + 1;
  END WHILE;
END;
```

Many system variables are dynamic and can be changed at runtime by using the SET statement. For a list, see Section 5.1.6.2, “Dynamic System Variables”. To change a system variable with SET, refer to it by name, optionally preceded by a modifier:

- To indicate that a variable is a global variable, precede its name by the GLOBAL keyword or the @@global. qualifier:

  ```
  SET GLOBAL max_connections = 1000;
  SET @@global.max_connections = 1000;
  ```

  The SYSTEM_VARIABLES_ADMIN or SUPER privilege is required to set global variables.

- Another way to set a global variable is to precede its name by the PERSIST keyword or the @@persist. qualifier:

  ```
  SET PERSIST max_connections = 1000;
  SET @@persist.max_connections = 1000;
  ```

  This SET syntax enables you to make configuration changes at runtime that also persist across server restarts. Like SET GLOBAL, SET PERSIST changes the runtime variable value, but also writes the variable setting to an option file named mysqld-auto.cnf in the data directory (replacing any existing variable setting if there is one). At startup, the server processes this file after all other

option files. The SYSTEM_VARIABLES_ADMIN or SUPER privilege is required to persist global variables.

Note

Management of the mysqld-auto.cnf file should be left to the server and not performed manually:

- Removal of the file results in a loss of all persisted settings at the next server startup. (This is permissible if your intent is to reconfigure the server without these settings.) To remove all settings in the file without removing the file itself, use this statement:

```
RESET PERSIST;
```

- Manual changes to the file may result in a parse error at server startup. In this case, the server reports an error and exits. If this issue occurs, start the server with the persisted_globals_load system variable disabled or with the --no-defaults option. Alternatively, remove the mysqld-auto.cnf file, but, as noted previously, removing this file results in a loss of all persisted settings.

A plugin variable can be persisted if the plugin is installed when SET PERSIST is executed. Assignment of the persisted plugin variable takes effect for subsequent server restarts if the plugin is still installed. If the plugin is no longer installed, the plugin variable will not exist when the server reads the mysqld-auto.cnf file. In this case, the server writes a warning to the error log and continues:

```
currently unknown variable 'var_name'
was read from the persisted config file
```

- To indicate that a variable is a session variable, precede its name by the SESSION keyword or either the @@session. or @@ qualifier:

```
SET SESSION sql_mode = 'TRADITIONAL';
SET @@session.sql_mode = 'TRADITIONAL';
SET @@sql_mode = 'TRADITIONAL';
```

Setting a session variable normally requires no special privilege, although there are exceptions that require the SYSTEM_VARIABLES_ADMIN or SUPER privilege (such as sql_log_bin). A client can change its own session variables, but not those of any other client.

Session-only system variables cannot be persisted. They cannot be set at server startup, so there is no reason to list them in mysqld-auto.cnf.

- LOCAL and @@local. are synonyms for SESSION and @@session..
- If no modifier is present, SET changes the session variable.
- An error occurs under these circumstances:
  - Use of SET GLOBAL (or @@global.) or SET PERSIST (or @@persist.) when setting a variable that has only a session value:

```
mysql> SET GLOBAL sql_log_bin = ON;
ERROR 1231 (42000): Variable 'sql_log_bin' can't be
set to the value of 'ON'
```

  - Omission of GLOBAL (or @@global.) or PERSIST (or @@persist.) when setting a variable that has only a global value:

```
mysql> SET max_connections = 1000;
ERROR 1229 (HY000): Variable 'max_connections' is a
GLOBAL variable and should be set with SET GLOBAL
```

- Use of `SET SESSION` (or `@@SESSION.`) when setting a variable that has only a global value:

```
mysql> SET SESSION max_connections = 1000;
ERROR 1229 (HY000): Variable 'max_connections' is a
GLOBAL variable and should be set with SET GLOBAL
```

The preceding modifiers apply only to system variables. An error occurs for attempts to apply them to user-defined variables, stored procedure or function parameters, or stored program local variables.

A `SET` statement can contain multiple variable assignments, separated by commas. This statement assigns values to a user-defined variable and a system variable:

```
SET @x = 1, SESSION sql_mode = '';
```

If you set multiple system variables, the most recent `GLOBAL`, `PERSIST`, or `SESSION` modifier in the statement is used for following assignments that have no modifier specified.

Examples of multiple-variable assignment:

```
SET GLOBAL sort_buffer_size = 1000000, SESSION sort_buffer_size = 1000000;
SET @@global.sort_buffer_size = 1000000, @@local.sort_buffer_size = 1000000;
SET GLOBAL max_connections = 1000, sort_buffer_size = 1000000;
```

If any variable assignment in a `SET` statement fails, the entire statement fails and no variables are changed, nor is the `mysqld-auto.cnf` file changed.

If you change a session system variable, the value remains in effect within your session until you change the variable to a different value or the session ends. The change has no effect on other sessions.

If you change a global system variable, the value is remembered and used for new sessions until you change the variable to a different value or the server exits. The change is visible to any client that accesses the global variable. However, the change affects the corresponding session variable only for clients that connect after the change. The global variable change does not affect the session variable for any current client sessions (not even the session within which the `SET GLOBAL` statement occurred).

To make a global system variable setting permanent so that it applies across server restarts, modify it with `SET PERSIST` to record it in the `mysqld-auto.cnf` file. It is also possible to use `SET GLOBAL` and manually modify a `my.cnf` file, but that is more cumbersome, and an error in a manually entered setting might not be discovered until much later. `SET PERSIST` is more convenient and avoids the possibility of malformed settings.

The Performance Schema `persisted_variables` table provides an SQL interface to the `mysqld-auto.cnf` file, enabling its contents to be inspected at runtime using `SELECT` statements. See Section 25.11.13.1, “Performance Schema persisted_variables Table”.

The Performance Schema `variables_info` table contains information showing when and by which user each system variable was most recently set. See Section 25.11.13.2, “Performance Schema variables_info Table”.

To set a `GLOBAL` value to the compiled-in MySQL default value or a `SESSION` variable to the current corresponding `GLOBAL` value, set the variable to the value `DEFAULT`. For example, the following two statements are identical in setting the session value of `max_join_size` to the current global value:

```
SET @@session.max_join_size=DEFAULT;
SET @@session.max_join_size=@@global.max_join_size;
```

Not all system variables can be set to `DEFAULT`. In such cases, assigning `DEFAULT` results in an error.

With `SET PERSIST` (or `@@persist.`), setting a global variable to `DEFAULT` or to its literal default value assigns the variable its default value and adds a setting for the variable to `mysqld-auto.cnf`. To remove the variable from the file, use `RESET PERSIST`.

An error occurs for attempts to assign `DEFAULT` to user-defined variables, stored procedure or function parameters, or stored program local variables.

To refer to the value of a system variable in expressions, use one of the `@@`-modifiers (except `@@persist.`, which is not permitted in expressions). For example, you can retrieve values in a `SELECT` statement like this:

```
SELECT @@global.sql_mode, @@session.sql_mode, @@sql_mode;
```

For a reference to a system variable in an expression as *`@@var_name`* (rather than with `@@global.` or `@@session.`), MySQL returns the session value if it exists and the global value otherwise. This differs from `SET @@var_name = expr`, which always refers to the session value.

13.7.4.2 SET CHARACTER SET Syntax

```
SET {CHARACTER SET | CHARSET}
    {'charset_name' | DEFAULT}
```

This statement maps all strings sent between the server and the current client with the given mapping. `SET CHARACTER SET` sets three session system variables: `character_set_client` and `character_set_results` are set to the given character set, and `character_set_connection` to the value of `character_set_database`. See Section 10.1.4, “Connection Character Sets and Collations”.

`charset_name` may be quoted or unquoted.

The default character set mapping can be restored by using the value `DEFAULT`. The default depends on the server configuration.

`ucs2`, `utf16`, and `utf32` cannot be used as a client character set, which means that they do not work for `SET CHARACTER SET`.

13.7.4.3 SET NAMES Syntax

```
SET NAMES {'charset_name'
    [COLLATE 'collation_name'] | DEFAULT}
```

This statement sets the three session system variables `character_set_client`, `character_set_connection`, and `character_set_results` to the given character set. Setting `character_set_connection` to `charset_name` also sets `collation_connection` to the default collation for `charset_name`. See Section 10.1.4, “Connection Character Sets and Collations”.

The optional `COLLATE` clause may be used to specify a collation explicitly. If given, the collation must one of the permitted collations for *`charset_name`*.

`charset_name` and *`collation_name`* may be quoted or unquoted.

The default mapping can be restored by using a value of `DEFAULT`. The default depends on the server configuration.

`ucs2`, `utf16`, and `utf32` cannot be used as a client character set, which means that they do not work for `SET NAMES`.

13.7.5 SHOW Syntax

`SHOW` has many forms that provide information about databases, tables, columns, or status information about the server. This section describes those following:

```
SHOW {BINARY | MASTER} LOGS
SHOW BINLOG EVENTS [IN 'log_name'] [FROM pos] [LIMIT [offset,] row_count]
SHOW CHARACTER SET [like_or_where]
SHOW COLLATION [like_or_where]
SHOW [FULL] COLUMNS FROM tbl_name [FROM db_name] [like_or_where]
SHOW CREATE DATABASE db_name
SHOW CREATE EVENT event_name
SHOW CREATE FUNCTION func_name
SHOW CREATE PROCEDURE proc_name
SHOW CREATE TABLE tbl_name
SHOW CREATE TRIGGER trigger_name
SHOW CREATE VIEW view_name
SHOW DATABASES [like_or_where]
SHOW ENGINE engine_name {STATUS | MUTEX}
SHOW [STORAGE] ENGINES
SHOW ERRORS [LIMIT [offset,] row_count]
SHOW EVENTS
SHOW FUNCTION CODE func_name
SHOW FUNCTION STATUS [like_or_where]
SHOW GRANTS FOR user
SHOW INDEX FROM tbl_name [FROM db_name]
SHOW MASTER STATUS
SHOW OPEN TABLES [FROM db_name] [like_or_where]
SHOW PLUGINS
SHOW PROCEDURE CODE proc_name
SHOW PROCEDURE STATUS [like_or_where]
SHOW PRIVILEGES
SHOW [FULL] PROCESSLIST
SHOW PROFILE [types] [FOR QUERY n] [OFFSET n] [LIMIT n]
SHOW PROFILES
SHOW RELAYLOG EVENTS [IN 'log_name'] [FROM pos] [LIMIT [offset,] row_count]
SHOW SLAVE HOSTS
SHOW SLAVE STATUS [FOR CHANNEL channel]
SHOW [GLOBAL | SESSION] STATUS [like_or_where]
SHOW TABLE STATUS [FROM db_name] [like_or_where]
SHOW [FULL] TABLES [FROM db_name] [like_or_where]
SHOW TRIGGERS [FROM db_name] [like_or_where]
SHOW [GLOBAL | SESSION] VARIABLES [like_or_where]
SHOW WARNINGS [LIMIT [offset,] row_count]

like_or_where:
    LIKE 'pattern'
  | WHERE expr
```

If the syntax for a given `SHOW` statement includes a `LIKE 'pattern'` part, `'pattern'` is a string that can contain the SQL `%` and `_` wildcard characters. The pattern is useful for restricting statement output to matching values.

Several `SHOW` statements also accept a `WHERE` clause that provides more flexibility in specifying which rows to display. See Section 24.33, “Extensions to SHOW Statements”.

Many MySQL APIs (such as PHP) enable you to treat the result returned from a `SHOW` statement as you would a result set from a `SELECT`; see Chapter 27, *Connectors and APIs*, or your API documentation for more information. In addition, you can work in SQL with results from queries on tables in the `INFORMATION_SCHEMA` database, which you cannot easily do with results from `SHOW` statements. See Chapter 24, *INFORMATION_SCHEMA Tables*.

13.7.5.1 SHOW BINARY LOGS Syntax

```
SHOW BINARY LOGS
SHOW MASTER LOGS
```

Lists the binary log files on the server. This statement is used as part of the procedure described in Section 13.4.1.1, “PURGE BINARY LOGS Syntax”, that shows how to determine which logs can be purged.

```
mysql> SHOW BINARY LOGS;
+---------------+-----------+
| Log_name      | File_size |
+---------------+-----------+
| binlog.000015 |    724935 |
| binlog.000016 |    733481 |
+---------------+-----------+
```

`SHOW MASTER LOGS` is equivalent to `SHOW BINARY LOGS`.

A user with the `SUPER` or `REPLICATION CLIENT` privilege may execute this statement.

13.7.5.2 SHOW BINLOG EVENTS Syntax

```
SHOW BINLOG EVENTS
   [IN 'log_name']
   [FROM pos]
   [LIMIT [offset,] row_count]
```

Shows the events in the binary log. If you do not specify `'log_name'`, the first binary log is displayed.

The `LIMIT` clause has the same syntax as for the `SELECT` statement. See Section 13.2.10, “SELECT Syntax”.

Note

Issuing a `SHOW BINLOG EVENTS` with no `LIMIT` clause could start a very time- and resource-consuming process because the server returns to the client the complete contents of the binary log (which includes all statements executed by the server that modify data). As an alternative to `SHOW BINLOG EVENTS`, use the `mysqlbinlog` utility to save the binary log to a text file for later examination and analysis. See Section 4.6.8, “`mysqlbinlog` — Utility for Processing Binary Log Files”.

Note

Some events relating to the setting of user and system variables are not included in the output from `SHOW BINLOG EVENTS`. To get complete coverage of events within a binary log, use `mysqlbinlog`.

Note

`SHOW BINLOG EVENTS` does *not* work with relay log files. You can use `SHOW RELAYLOG EVENTS` for this purpose.

13.7.5.3 SHOW CHARACTER SET Syntax

```
SHOW CHARACTER SET
    [LIKE 'pattern' | WHERE expr]
```

The `SHOW CHARACTER SET` statement shows all available character sets. The `LIKE` clause, if present, indicates which character set names to match. The `WHERE` clause can be given to select rows using

more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”. For example:

```
mysql> SHOW CHARACTER SET LIKE 'latin%';
+---------+-----------------------------+-------------------+--------+
| Charset | Description                 | Default collation | Maxlen |
+---------+-----------------------------+-------------------+--------+
| latin1  | cp1252 West European        | latin1_swedish_ci |      1 |
| latin2  | ISO 8859-2 Central European | latin2_general_ci |      1 |
| latin5  | ISO 8859-9 Turkish          | latin5_turkish_ci |      1 |
| latin7  | ISO 8859-13 Baltic          | latin7_general_ci |      1 |
+---------+-----------------------------+-------------------+--------+
```

The `Maxlen` column shows the maximum number of bytes required to store one character.

The `filename` character set is for internal use only; consequently, `SHOW CHARACTER SET` does not display it.

You can also obtain information about character sets from `INFORMATION_SCHEMA`, which contains a `CHARACTER_SETS` table. See Section 24.1, “The INFORMATION_SCHEMA CHARACTER_SETS Table”.

13.7.5.4 SHOW COLLATION Syntax

```
SHOW COLLATION
    [LIKE 'pattern' | WHERE expr]
```

This statement lists collations supported by the server. By default, the output from `SHOW COLLATION` includes all available collations. The `LIKE` clause, if present, indicates which collation names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”. For example:

```
mysql> SHOW COLLATION WHERE Charset = 'latin1';
+-------------------+---------+----+---------+----------+---------+
| Collation         | Charset | Id | Default | Compiled | Sortlen |
+-------------------+---------+----+---------+----------+---------+
| latin1_german1_ci | latin1  |  5 |         | Yes      |       1 |
| latin1_swedish_ci | latin1  |  8 | Yes     | Yes      |       1 |
| latin1_danish_ci  | latin1  | 15 |         | Yes      |       1 |
| latin1_german2_ci | latin1  | 31 |         | Yes      |       2 |
| latin1_bin        | latin1  | 47 |         | Yes      |       1 |
| latin1_general_ci | latin1  | 48 |         | Yes      |       1 |
| latin1_general_cs | latin1  | 49 |         | Yes      |       1 |
| latin1_spanish_ci | latin1  | 94 |         | Yes      |       1 |
+-------------------+---------+----+---------+----------+---------+
```

The `Collation` and `Charset` columns indicate the names of the collation and the character set with which it is associated. `Id` is the collation ID. `Default` indicates whether the collation is the default for its character set. `Compiled` indicates whether the character set is compiled into the server. `Sortlen` is related to the amount of memory required to sort strings expressed in the character set.

To see the default collation for each character set, use the following statement. `Default` is a reserved word, so to use it as an identifier, it must be quoted as such:

```
mysql> SHOW COLLATION WHERE `Default` = 'Yes';
+--------------------+----------+----+---------+----------+---------+
| Collation          | Charset  | Id | Default | Compiled | Sortlen |
+--------------------+----------+----+---------+----------+---------+
| big5_chinese_ci    | big5     |  1 | Yes     | Yes      |       1 |
| dec8_swedish_ci    | dec8     |  3 | Yes     | Yes      |       1 |
| cp850_general_ci   | cp850    |  4 | Yes     | Yes      |       1 |
| hp8_english_ci     | hp8      |  6 | Yes     | Yes      |       1 |
| koi8r_general_ci   | koi8r    |  7 | Yes     | Yes      |       1 |
```

```
| latin1_swedish_ci   | latin1   |  8 | Yes     | Yes      |       1 |
...
```

You can also obtain information about collations from `INFORMATION_SCHEMA`, which contains a `COLLATIONS` table. See Section 24.2, “The INFORMATION_SCHEMA COLLATIONS Table”.

13.7.5.5 SHOW COLUMNS Syntax

```
SHOW [EXTENDED] [FULL] {COLUMNS | FIELDS}
    {FROM | IN} tbl_name
    [{FROM | IN} db_name]
    [LIKE 'pattern' | WHERE expr]
```

`SHOW COLUMNS` displays information about the columns in a given table. It also works for views. `SHOW COLUMNS` displays information only for those columns for which you have some privilege.

You can use *`db_name.tbl_name`* as an alternative to the *`tbl_name FROM db_name`* syntax. In other words, these two statements are equivalent:

```
SHOW COLUMNS FROM mytable FROM mydb;
SHOW COLUMNS FROM mydb.mytable;
```

As of MySQL 8.0.2, the optional `EXTENDED` keyword causes the output to include information about hidden columns that MySQL uses internally and are not accessible by users.

The optional `FULL` keyword causes the output to include the column collation and comments, as well as the privileges you have for each column.

The `LIKE` clause, if present, indicates which column names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

```
mysql> SHOW COLUMNS FROM City;
+-------------+----------+------+-----+---------+----------------+
| Field       | Type     | Null | Key | Default | Extra          |
+-------------+----------+------+-----+---------+----------------+
| ID          | int(11)  | NO   | PRI | NULL    | auto_increment |
| Name        | char(35) | NO   |     |         |                |
| CountryCode | char(3)  | NO   | MUL |         |                |
| District    | char(20) | NO   |     |         |                |
| Population  | int(11)  | NO   |     | 0       |                |
+-------------+----------+------+-----+---------+----------------+
```

The data types may differ from what you expect them to be based on a `CREATE TABLE` statement because MySQL sometimes changes data types when you create or alter a table. The conditions under which this occurs are described in Section 13.1.15.7, “Silent Column Specification Changes”.

`SHOW COLUMNS` displays the following values for each table column:

- `Field`

  The column name.

- `Type`

  The column data type.

- `Collation`

  The collation for nonbinary string columns, or `NULL` for other columns. This value is displayed only if you use the `FULL` keyword.

- `Null`

Column nullability. The value is `YES` if `NULL` values can be stored in the column, `NO` if not.

- `Key`

  Whether the column is indexed:

  - If `Key` is empty, the column either is not indexed or is indexed only as a secondary column in a multiple-column, nonunique index.
  - If `Key` is `PRI`, the column is a `PRIMARY KEY` or is one of the columns in a multiple-column `PRIMARY KEY`.
  - If `Key` is `UNI`, the column is the first column of a `UNIQUE` index. (A `UNIQUE` index permits multiple `NULL` values, but you can tell whether the column permits `NULL` by checking the `Null` field.)
  - If `Key` is `MUL`, the column is the first column of a nonunique index in which multiple occurrences of a given value are permitted within the column.

  If more than one of the `Key` values applies to a given column of a table, `Key` displays the one with the highest priority, in the order `PRI`, `UNI`, `MUL`.

  A `UNIQUE` index may be displayed as `PRI` if it cannot contain `NULL` values and there is no `PRIMARY KEY` in the table. A `UNIQUE` index may display as `MUL` if several columns form a composite `UNIQUE` index; although the combination of the columns is unique, each column can still hold multiple occurrences of a given value.

- `Default`

  The default value for the column. This is `NULL` if the column has an explicit default of `NULL`, or if the column definition includes no `DEFAULT` clause.

- `Extra`

  Any additional information that is available about a given column. The value is nonempty in these cases:

  - `auto_increment` for columns that have the `AUTO_INCREMENT` attribute
  - `on update CURRENT_TIMESTAMP` for `TIMESTAMP` or `DATETIME` columns that have the `ON UPDATE CURRENT_TIMESTAMP` attribute
  - `VIRTUAL GENERATED` or `VIRTUAL STORED` for generated columns

- `Privileges`

  The privileges you have for the column. This value is displayed only if you use the `FULL` keyword.

- `Comment`

  Any comment included in the column definition. This value is displayed only if you use the `FULL` keyword.

You can also obtain information about table columns from `INFORMATION_SCHEMA`, which contains a `COLUMNS` table. See Section 24.4, “The INFORMATION_SCHEMA COLUMNS Table”. The extended information about hidden columns is available only using `SHOW EXTENDED COLUMNS`; it cannot be obtained from the `COLUMNS` table.

You can list a table's columns with the `mysqlshow *db_name tbl_name*` command.

The `DESCRIBE` statement provides information similar to `SHOW COLUMNS`. See Section 13.8.1, “DESCRIBE Syntax”.

The SHOW CREATE TABLE, SHOW TABLE STATUS, and SHOW INDEX statements also provide information about tables. See Section 13.7.5, “SHOW Syntax”.

13.7.5.6 SHOW CREATE DATABASE Syntax

```
SHOW CREATE {DATABASE | SCHEMA} [IF NOT EXISTS] db_name
```

Shows the CREATE DATABASE statement that creates the named database. If the SHOW statement includes an IF NOT EXISTS clause, the output too includes such a clause. SHOW CREATE SCHEMA is a synonym for SHOW CREATE DATABASE.

```
mysql> SHOW CREATE DATABASE test\G
*************************** 1. row ***************************
       Database: test
Create Database: CREATE DATABASE `test`
                 /*!40100 DEFAULT CHARACTER SET utf8mb4 */

mysql> SHOW CREATE SCHEMA test\G
*************************** 1. row ***************************
       Database: test
Create Database: CREATE DATABASE `test`
                 /*!40100 DEFAULT CHARACTER SET utf8mb4 */
```

SHOW CREATE DATABASE quotes table and column names according to the value of the sql_quote_show_create option. See Section 5.1.5, “Server System Variables”.

13.7.5.7 SHOW CREATE EVENT Syntax

```
SHOW CREATE EVENT event_name
```

This statement displays the CREATE EVENT statement needed to re-create a given event. It requires the EVENT privilege for the database from which the event is to be shown. For example (using the same event e_daily defined and then altered in Section 13.7.5.18, “SHOW EVENTS Syntax”):

```
mysql> SHOW CREATE EVENT test.e_daily\G
*************************** 1. row ***************************
               Event: e_daily
            sql_mode:
           time_zone: SYSTEM
        Create Event: CREATE EVENT `e_daily`
                        ON SCHEDULE EVERY 1 DAY
                        STARTS CURRENT_TIMESTAMP + INTERVAL 6 HOUR
                        ON COMPLETION NOT PRESERVE
                        ENABLE
                        COMMENT 'Saves total number of sessions then
                                clears the table each day'
                        DO BEGIN
                          INSERT INTO site_activity.totals (time, total)
                            SELECT CURRENT_TIMESTAMP, COUNT(*)
                            FROM site_activity.sessions;
                          DELETE FROM site_activity.sessions;
                        END
character_set_client: utf8
collation_connection: utf8_general_ci
  Database Collation: utf8mb4_0900_ai_ci
```

character_set_client is the session value of the character_set_client system variable when the event was created. collation_connection is the session value of the collation_connection system variable when the event was created. Database Collation is the collation of the database with which the event is associated.

The output reflects the current status of the event (ENABLE) rather than the status with which it was created.

13.7.5.8 SHOW CREATE FUNCTION Syntax

```
SHOW CREATE FUNCTION func_name
```

This statement is similar to `SHOW CREATE PROCEDURE` but for stored functions. See Section 13.7.5.9, “SHOW CREATE PROCEDURE Syntax”.

13.7.5.9 SHOW CREATE PROCEDURE Syntax

```
SHOW CREATE PROCEDURE proc_name
```

This statement is a MySQL extension. It returns the exact string that can be used to re-create the named stored procedure. A similar statement, `SHOW CREATE FUNCTION`, displays information about stored functions (see Section 13.7.5.8, “SHOW CREATE FUNCTION Syntax”).

To use either statement, you must have the global `SELECT` privilege.

```
mysql> SHOW CREATE PROCEDURE test.simpleproc\G
*************************** 1. row ***************************
           Procedure: simpleproc
            sql_mode:
    Create Procedure: CREATE PROCEDURE `simpleproc`(OUT param1 INT)
                      BEGIN
                      SELECT COUNT(*) INTO param1 FROM t;
                      END
character_set_client: utf8
collation_connection: utf8_general_ci
  Database Collation: utf8mb4_0900_ai_ci

mysql> SHOW CREATE FUNCTION test.hello\G
*************************** 1. row ***************************
            Function: hello
            sql_mode:
     Create Function: CREATE FUNCTION `hello`(s CHAR(20))
                      RETURNS CHAR(50)
                      RETURN CONCAT('Hello, ',s,'!')
character_set_client: utf8
collation_connection: utf8_general_ci
  Database Collation: utf8mb4_0900_ai_ci
```

`character_set_client` is the session value of the `character_set_client` system variable when the routine was created. `collation_connection` is the session value of the `collation_connection` system variable when the routine was created. `Database Collation` is the collation of the database with which the routine is associated.

13.7.5.10 SHOW CREATE TABLE Syntax

```
SHOW CREATE TABLE tbl_name
```

Shows the `CREATE TABLE` statement that creates the named table. To use this statement, you must have some privilege for the table. This statement also works with views.

```
mysql> SHOW CREATE TABLE t\G
*************************** 1. row ***************************
       Table: t
Create Table: CREATE TABLE `t` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `s` char(60) DEFAULT NULL,
  PRIMARY KEY (`id`)
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4
```

`SHOW CREATE TABLE` quotes table and column names according to the value of the `sql_quote_show_create` option. See Section 5.1.5, “Server System Variables”.

13.7.5.11 SHOW CREATE TRIGGER Syntax

```
SHOW CREATE TRIGGER trigger_name
```

This statement shows the `CREATE TRIGGER` statement that creates the named trigger. This statement requires the `TRIGGER` privilege for the table associated with the trigger.

```
mysql> SHOW CREATE TRIGGER ins_sum\G
*************************** 1. row ***************************
               Trigger: ins_sum
              sql_mode: STRICT_TRANS_TABLES,NO_ENGINE_SUBSTITUTION
SQL Original Statement: CREATE DEFINER=`me`@`localhost` TRIGGER ins_sum
                        BEFORE INSERT ON account
                        FOR EACH ROW SET @sum = @sum + NEW.amount
  character_set_client: utf8
  collation_connection: utf8_general_ci
    Database Collation: utf8mb4_0900_ai_ci
               Created: 2013-07-09 10:39:34.96
```

`SHOW CREATE TRIGGER` output has the following columns:

- `Trigger`: The trigger name.
- `sql_mode`: The SQL mode in effect when the trigger executes.
- `SQL Original Statement`: The `CREATE TRIGGER` statement that defines the trigger.
- `character_set_client`: The session value of the `character_set_client` system variable when the trigger was created.
- `collation_connection`: The session value of the `collation_connection` system variable when the trigger was created.
- `Database Collation`: The collation of the database with which the trigger is associated.
- `Created`: The date and time when the trigger was created. This is a `TIMESTAMP(2)` value (with a fractional part in hundredths of seconds) for triggers.

You can also obtain information about trigger objects from `INFORMATION_SCHEMA`, which contains a `TRIGGERS` table. See Section 24.28, “The INFORMATION_SCHEMA TRIGGERS Table”.

13.7.5.12 SHOW CREATE USER Syntax

```
SHOW CREATE USER user
```

This statement shows the `CREATE USER` statement that creates the named user. An error occurs if the user does not exist. The statement requires the `SELECT` privilege for the `mysql` database, except to see the privileges for the current user.

To name the account, use the format described in Section 6.2.4, “Specifying Account Names”. The host name part of the account name, if omitted, defaults to `'%'`. It is also possible to specify `CURRENT_USER` or `CURRENT_USER()` to refer to the account associated with the current session.

```
mysql> SHOW CREATE USER 'root'@'localhost'\G
*************************** 1. row ***************************
CREATE USER for root@localhost: CREATE USER 'root'@'localhost'
IDENTIFIED WITH 'mysql_native_password'
AS '*2470C0C06DEE42FD1618BB99005ADCA2EC9D1E19'
REQUIRE NONE PASSWORD EXPIRE DEFAULT ACCOUNT UNLOCK
```

The output format is affected by the setting of the `log_builtin_as_identified_by_password` system variable.

To display the privileges granted to an account, use the `SHOW GRANTS` statement. See Section 13.7.5.21, “SHOW GRANTS Syntax”.

13.7.5.13 SHOW CREATE VIEW Syntax

```
SHOW CREATE VIEW view_name
```

This statement shows the `CREATE VIEW` statement that creates the named view.

```
mysql> SHOW CREATE VIEW v\G
*************************** 1. row ***************************
                View: v
         Create View: CREATE ALGORITHM=UNDEFINED
                      DEFINER=`bob`@`localhost`
                      SQL SECURITY DEFINER VIEW
                      `v` AS select 1 AS `a`,2 AS `b`
character_set_client: utf8
collation_connection: utf8_general_ci
```

`character_set_client` is the session value of the `character_set_client` system variable when the view was created. `collation_connection` is the session value of the `collation_connection` system variable when the view was created.

Use of `SHOW CREATE VIEW` requires the `SHOW VIEW` privilege, and the `SELECT` privilege for the view in question.

You can also obtain information about view objects from `INFORMATION_SCHEMA`, which contains a `VIEWS` table. See Section 24.30, “The INFORMATION_SCHEMA VIEWS Table”.

MySQL lets you use different `sql_mode` settings to tell the server the type of SQL syntax to support. For example, you might use the `ANSI` SQL mode to ensure MySQL correctly interprets the standard SQL concatenation operator, the double bar (`||`), in your queries. If you then create a view that concatenates items, you might worry that changing the `sql_mode` setting to a value different from `ANSI` could cause the view to become invalid. But this is not the case. No matter how you write out a view definition, MySQL always stores it the same way, in a canonical form. Here is an example that shows how the server changes a double bar concatenation operator to a `CONCAT()` function:

```
mysql> SET sql_mode = 'ANSI';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE VIEW test.v AS SELECT 'a' || 'b' as col1;
Query OK, 0 rows affected (0.01 sec)

mysql> SHOW CREATE VIEW test.v\G
*************************** 1. row ***************************
                View: v
         Create View: CREATE VIEW "v" AS select concat('a','b') AS "col1"
...
1 row in set (0.00 sec)
```

The advantage of storing a view definition in canonical form is that changes made later to the value of `sql_mode` will not affect the results from the view. However an additional consequence is that comments prior to `SELECT` are stripped from the definition by the server.

13.7.5.14 SHOW DATABASES Syntax

```
SHOW {DATABASES | SCHEMAS}
    [LIKE 'pattern' | WHERE expr]
```

`SHOW DATABASES` lists the databases on the MySQL server host. `SHOW SCHEMAS` is a synonym for `SHOW DATABASES`. The `LIKE` clause, if present, indicates which database names to match.

The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

You see only those databases for which you have some kind of privilege, unless you have the global `SHOW DATABASES` privilege. You can also get this list using the `mysqlshow` command.

If the server was started with the `--skip-show-database` option, you cannot use this statement at all unless you have the `SHOW DATABASES` privilege.

MySQL implements databases as directories in the data directory, so this statement simply lists directories in that location. However, the output may include names of directories that do not correspond to actual databases.

13.7.5.15 SHOW ENGINE Syntax

```
SHOW ENGINE engine_name {STATUS | MUTEX}
```

`SHOW ENGINE` displays operational information about a storage engine. It requires the `PROCESS` privilege. The statement has these variants:

```
SHOW ENGINE INNODB STATUS
SHOW ENGINE INNODB MUTEX
SHOW ENGINE PERFORMANCE_SCHEMA STATUS
```

`SHOW ENGINE INNODB STATUS` displays extensive information from the standard `InnoDB` Monitor about the state of the `InnoDB` storage engine. For information about the standard monitor and other `InnoDB` Monitors that provide information about `InnoDB` processing, see Section 15.16, “InnoDB Monitors”.

`SHOW ENGINE INNODB MUTEX` displays `InnoDB` mutex and rw-lock statistics.

Note

`InnoDB` mutexes and rwlocks can also be monitored using Performance Schema tables. See Section 15.15.2, “Monitoring InnoDB Mutex Waits Using Performance Schema”.

Mutex statistics collection is configured dynamically using the following options:

- To enable the collection of mutex statistics, run:

  ```
  SET GLOBAL innodb_monitor_enable='latch';
  ```

- To reset mutex statistics, run:

  ```
  SET GLOBAL innodb_monitor_reset='latch';
  ```

- To disable the collection of mutex statistics, run:

  ```
  SET GLOBAL innodb_monitor_disable='latch';
  ```

Collection of mutex statistics for `SHOW ENGINE INNODB MUTEX` can also be enabled by setting `innodb_monitor_enable='all'`, or disabled by setting `innodb_monitor_disable='all'`.

`SHOW ENGINE INNODB MUTEX` output has the following columns:

- `Type`

  Always `InnoDB`.

- `Name`

For mutexes, the `Name` field reports only the mutex name. For rwlocks, the `Name` field reports the source file where the rwlock is implemented, and the line number in the file where the rwlock is created. The line number is specific to your version of MySQL.

- `Status`

  The mutex status. This field reports the number of spins, waits, and calls. Statistics for low-level operating system mutexes, which are implemented outside of `InnoDB`, are not reported.

  - `spins` indicates the number of spins.
  - `waits` indicates the number of mutex waits.
  - `calls` indicates how many times the mutex was requested.

`SHOW ENGINE INNODB MUTEX` skips the mutexes and rw-locks of buffer pool blocks, as the amount of output can be overwhelming on systems with a large buffer pool. (There is one mutex and one rw-lock in each 16K buffer pool block, and there are 65,536 blocks per gigabyte.) `SHOW ENGINE INNODB MUTEX` also does not list any mutexes or rw-locks that have never been waited on (`os_waits=0`). Thus, `SHOW ENGINE INNODB MUTEX` only displays information about mutexes and rw-locks outside of the buffer pool that have caused at least one OS-level wait.

Use `SHOW ENGINE PERFORMANCE_SCHEMA STATUS` to inspect the internal operation of the Performance Schema code:

```
mysql> SHOW ENGINE PERFORMANCE_SCHEMA STATUS\G
...
*************************** 3. row ***************************
  Type: performance_schema
  Name: events_waits_history.size
Status: 76
*************************** 4. row ***************************
  Type: performance_schema
  Name: events_waits_history.count
Status: 10000
*************************** 5. row ***************************
  Type: performance_schema
  Name: events_waits_history.memory
Status: 760000
...
*************************** 57. row ***************************
  Type: performance_schema
  Name: performance_schema.memory
Status: 26459600
...
```

This statement is intended to help the DBA understand the effects that different Performance Schema options have on memory requirements.

`Name` values consist of two parts, which name an internal buffer and a buffer attribute, respectively. Interpret buffer names as follows:

- An internal buffer that is not exposed as a table is named within parentheses. Examples: `(pfs_cond_class).size`, `(pfs_mutex_class).memory`.
- An internal buffer that is exposed as a table in the `performance_schema` database is named after the table, without parentheses. Examples: `events_waits_history.size`, `mutex_instances.count`.
- A value that applies to the Performance Schema as a whole begins with `performance_schema`. Example: `performance_schema.memory`.

Buffer attributes have these meanings:

- `size` is the size of the internal record used by the implementation, such as the size of a row in a table. `size` values cannot be changed.
- `count` is the number of internal records, such as the number of rows in a table. `count` values can be changed using Performance Schema configuration options.
- For a table, *`tbl_name`*`.memory` is the product of `size` and `count`. For the Performance Schema as a whole, `performance_schema.memory` is the sum of all the memory used (the sum of all other `memory` values).

In some cases, there is a direct relationship between a Performance Schema configuration parameter and a `SHOW ENGINE` value. For example, `events_waits_history_long.count` corresponds to `performance_schema_events_waits_history_long_size`. In other cases, the relationship is more complex. For example, `events_waits_history.count` corresponds to `performance_schema_events_waits_history_size` (the number of rows per thread) multiplied by `performance_schema_max_thread_instances` (the number of threads).

13.7.5.16 SHOW ENGINES Syntax

```
SHOW [STORAGE] ENGINES
```

`SHOW ENGINES` displays status information about the server's storage engines. This is particularly useful for checking whether a storage engine is supported, or to see what the default engine is. This information can also be obtained from the `INFORMATION_SCHEMA ENGINES` table. See Section 24.7, “The INFORMATION_SCHEMA ENGINES Table”.

```
mysql> SHOW ENGINES\G
*************************** 1. row ***************************
      Engine: InnoDB
     Support: DEFAULT
     Comment: Supports transactions, row-level locking, and foreign keys
Transactions: YES
          XA: YES
  Savepoints: YES
*************************** 2. row ***************************
      Engine: MRG_MYISAM
     Support: YES
     Comment: Collection of identical MyISAM tables
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 3. row ***************************
      Engine: MEMORY
     Support: YES
     Comment: Hash based, stored in memory, useful for temporary tables
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 4. row ***************************
      Engine: BLACKHOLE
     Support: YES
     Comment: /dev/null storage engine (anything you write to it disappears)
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 5. row ***************************
      Engine: MyISAM
     Support: YES
     Comment: MyISAM storage engine
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 6. row ***************************
      Engine: CSV
     Support: YES
```

```
     Comment: CSV storage engine
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 7. row ***************************
      Engine: ARCHIVE
     Support: YES
     Comment: Archive storage engine
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 8. row ***************************
      Engine: PERFORMANCE_SCHEMA
     Support: YES
     Comment: Performance Schema
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 9. row ***************************
      Engine: FEDERATED
     Support: YES
     Comment: Federated MySQL storage engine
Transactions: NO
          XA: NO
  Savepoints: NO
```

The output from `SHOW ENGINES` may vary according to the MySQL version used and other factors. The values shown in the `Support` column indicate the server's level of support for the storage engine, as shown in the following table.

Value	Meaning
`YES`	The engine is supported and is active
`DEFAULT`	Like `YES`, plus this is the default engine
`NO`	The engine is not supported
`DISABLED`	The engine is supported but has been disabled

A value of `NO` means that the server was compiled without support for the engine, so it cannot be enabled at runtime.

A value of `DISABLED` occurs either because the server was started with an option that disables the engine, or because not all options required to enable it were given. In the latter case, the error log file should contain a reason indicating why the option is disabled. See Section 5.4.2, “The Error Log”.

You might also see `DISABLED` for a storage engine if the server was compiled to support it, but was started with a `--skip-`*`engine_name`* option.

All MySQL servers support `MyISAM` tables. It is not possible to disable `MyISAM`.

The `Transactions`, `XA`, and `Savepoints` columns indicate whether the storage engine supports transactions, XA transactions, and savepoints, respectively.

13.7.5.17 SHOW ERRORS Syntax

```
SHOW ERRORS [LIMIT [offset,] row_count]
SHOW COUNT(*) ERRORS
```

`SHOW ERRORS` is a diagnostic statement that is similar to `SHOW WARNINGS`, except that it displays information only for errors, rather than for errors, warnings, and notes.

The `LIMIT` clause has the same syntax as for the `SELECT` statement. See Section 13.2.10, “SELECT Syntax”.

The `SHOW COUNT(*) ERRORS` statement displays the number of errors. You can also retrieve this number from the `error_count` variable:

```
SHOW COUNT(*) ERRORS;
SELECT @@error_count;
```

`SHOW ERRORS` and `error_count` apply only to errors, not warnings or notes. In other respects, they are similar to `SHOW WARNINGS` and `warning_count`. In particular, `SHOW ERRORS` cannot display information for more than `max_error_count` messages, and `error_count` can exceed the value of `max_error_count` if the number of errors exceeds `max_error_count`.

For more information, see Section 13.7.5.40, “SHOW WARNINGS Syntax”.

13.7.5.18 SHOW EVENTS Syntax

```
SHOW EVENTS
    [{FROM | IN} schema_name]
    [LIKE 'pattern' | WHERE expr]
```

This statement displays information about Event Manager events. It requires the `EVENT` privilege for the database from which the events are to be shown.

In its simplest form, `SHOW EVENTS` lists all of the events in the current schema:

```
mysql> SELECT CURRENT_USER(), SCHEMA();
+----------------+----------+
| CURRENT_USER() | SCHEMA() |
+----------------+----------+
| jon@ghidora    | myschema |
+----------------+----------+
1 row in set (0.00 sec)

mysql> SHOW EVENTS\G
*************************** 1. row ***************************
                  Db: myschema
                Name: e_daily
             Definer: jon@ghidora
           Time zone: SYSTEM
                Type: RECURRING
          Execute at: NULL
      Interval value: 10
      Interval field: SECOND
              Starts: 2006-02-09 10:41:23
                Ends: NULL
              Status: ENABLED
          Originator: 0
character_set_client: utf8
collation_connection: utf8_general_ci
  Database Collation: utf8mb4_0900_ai_ci
```

To see events for a specific schema, use the `FROM` clause. For example, to see events for the `test` schema, use the following statement:

```
SHOW EVENTS FROM test;
```

The `LIKE` clause, if present, indicates which event names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

`SHOW EVENTS` output has the following columns:

- `Db`: The schema (database) on which the event is defined.

- `Name`: The name of the event.
- `Time zone`: The event time zone, which is the time zone used for scheduling the event and that is in effect within the event as it executes. The default value is `SYSTEM`.
- `Definer`: The account of the user who created the event, in `'user_name'@'host_name'` format.
- `Type`: The event repetition type, either `ONE TIME` (transient) or `RECURRING` (repeating).
- `Execute At`: The date and time when a transient event is set to execute. Shown as a `DATETIME` value.

  For a recurring event, the value of this column is always `NULL`.
- `Interval Value`: For a recurring event, the number of intervals to wait between event executions.

  For a transient event, the value of this column is always `NULL`.
- `Interval Field`: The time units used for the interval which a recurring event waits before repeating.

  For a transient event, the value of this column is always `NULL`.
- `Starts`: The start date and time for a recurring event. This is displayed as a `DATETIME` value, and is `NULL` if no start date and time are defined for the event.

  For a transient event, this column is always `NULL`.
- `Ends`: The end date and time for a recurring event. This is displayed as a `DATETIME` value, and defaults to `NULL` if no end date and time is defined for the event.

  For a transient event, this column is always `NULL`.
- `Status`: The event status. One of `ENABLED`, `DISABLED`, or `SLAVESIDE_DISABLED`.

  `SLAVESIDE_DISABLED` indicates that the creation of the event occurred on another MySQL server acting as a replication master and replicated to the current MySQL server which is acting as a slave, but the event is not presently being executed on the slave.
- `Originator`: The server ID of the MySQL server on which the event was created. Defaults to 0.
- `character_set_client` is the session value of the `character_set_client` system variable when the routine was created. `collation_connection` is the session value of the `collation_connection` system variable when the routine was created. `Database Collation` is the collation of the database with which the routine is associated.

For more information about `SLAVE_DISABLED` and the `Originator` column, see Section 18.4.1.12, “Replication of Invoked Features”.

The event action statement is not shown in the output of `SHOW EVENTS`. Use `SHOW CREATE EVENT` or the `INFORMATION_SCHEMA.EVENTS` table.

Times displayed by `SHOW EVENTS` are given in the event time zone, as discussed in Section 23.4.4, “Event Metadata”.

The columns in the output of `SHOW EVENTS` are similar to, but not identical to the columns in the `INFORMATION_SCHEMA.EVENTS` table. See Section 24.8, “The INFORMATION_SCHEMA EVENTS Table”.

13.7.5.19 SHOW FUNCTION CODE Syntax

```
SHOW FUNCTION CODE func_name
```

This statement is similar to `SHOW PROCEDURE CODE` but for stored functions. See Section 13.7.5.27, “SHOW PROCEDURE CODE Syntax”.

13.7.5.20 SHOW FUNCTION STATUS Syntax

```
SHOW FUNCTION STATUS
    [LIKE 'pattern' | WHERE expr]
```

This statement is similar to `SHOW PROCEDURE STATUS` but for stored functions. See Section 13.7.5.28, “SHOW PROCEDURE STATUS Syntax”.

13.7.5.21 SHOW GRANTS Syntax

```
SHOW GRANTS
    [FOR user_or_role
        [USING user [, user] ...]]

user_or_role: {
    user
  | role
}
```

This statement displays the privileges and roles that are assigned to a MySQL user account or role, in the form of `GRANT` statements that must be executed to duplicate the privilege and role assignments. `SHOW GRANTS` requires the `SELECT` privilege for the `mysql` database, except to see the privileges for the current user.

To name the account or role, use the same format as for the `GRANT` statement; for example, `'jeffrey'@'localhost'`. If you specify only the user name part of the account name, a host name part of `'%'` is used. For additional information about specifying account names, see Section 13.7.1.6, “GRANT Syntax”.

```
mysql> SHOW GRANTS FOR 'root'@'localhost';
+---------------------------------------------------------------------+
| Grants for root@localhost                                           |
+---------------------------------------------------------------------+
| GRANT ALL PRIVILEGES ON *.* TO 'root'@'localhost' WITH GRANT OPTION |
+---------------------------------------------------------------------+
```

In MySQL 8.0 compared to previous series, `SHOW GRANTS` no longer displays `ALL PRIVILEGES` in its global-privileges output because the meaning of `ALL PRIVILEGES` at the global level varies depending on which dynamic privileges are defined. Instead, `SHOW GRANTS` explictly lists each granted global privilege. Applications that process `SHOW GRANTS` output should be adjusted accordingly.

```
mysql> SHOW GRANTS FOR 'root'@'localhost';
+---------------------------------------------------------------------+
| Grants for root@localhost                                           |
+---------------------------------------------------------------------+
| GRANT SELECT, INSERT, UPDATE, DELETE, CREATE, DROP, RELOAD,         |
| SHUTDOWN, PROCESS, FILE, REFERENCES, INDEX, ALTER, SHOW DATABASES,  |
| SUPER, CREATE TEMPORARY TABLES, LOCK TABLES, EXECUTE, REPLICATION   |
| SLAVE, REPLICATION CLIENT, CREATE VIEW, SHOW VIEW, CREATE ROUTINE,  |
| ALTER ROUTINE, CREATE USER, EVENT, TRIGGER, CREATE TABLESPACE,      |
| CREATE ROLE, DROP ROLE ON *.* TO `root`@`localhost` WITH GRANT      |
| OPTION                                                              |
| GRANT PROXY ON ''@'' TO 'root'@'localhost' WITH GRANT OPTION        |
+---------------------------------------------------------------------+
```

At the global level, `GRANT OPTION` applies to all grant global privileges if granted for any of them, but applies individually to granted dynamic privileges. `SHOW GRANTS` displays global privileges this way:

- One line listing all granted static privileges, if there are any, including `WITH GRANT OPTION` if appropriate.
- One line listing all granted dynamic privileges for which `GRANT OPTION` is granted, if there are any, including `WITH GRANT OPTION`.
- One line listing all granted dynamic privileges for which `GRANT OPTION` is not granted, if there are any.

To display the privileges granted to the account that you are using to connect to the server, you can use any of the following statements:

```
SHOW GRANTS;
SHOW GRANTS FOR CURRENT_USER;
SHOW GRANTS FOR CURRENT_USER();
```

If `SHOW GRANTS FOR CURRENT_USER` (or any of the equivalent syntaxes) is used in `DEFINER` context, such as within a stored procedure that is defined with `SQL SECURITY DEFINER`), the grants displayed are those of the definer and not the invoker.

With the optional `USING` clause, `SHOW GRANTS` enables you to examine the privileges associated with roles for the user. Each role named in the `USING` clause must be granted to the user.

Unless you are displaying information for your own account, the `USING` clause requires the `SUPER` privilege.

Suppose that user `u1` is assigned roles `r1` and `r2`, as follows:

```
CREATE ROLE 'r1', 'r2';
GRANT SELECT ON db1.* TO 'r1';
GRANT INSERT, UPDATE, DELETE ON db1.* TO 'r2';
CREATE USER 'u1'@'localhost' IDENTIFIED BY 'u1pass';
GRANT 'r1', 'r2' TO 'u1'@'localhost';
```

`SHOW GRANTS` without `USING` shows the granted roles:

```
mysql> SHOW GRANTS FOR 'u1'@'localhost';
+-----------------------------------------+
| Grants for u1@localhost                 |
+-----------------------------------------+
| GRANT USAGE ON *.* TO `u1`@`localhost`  |
| GRANT `r1`@`%`,`r2`@`%` TO `u1`@`localhost` |
+-----------------------------------------+
```

Adding a `USING` clause causes the statement to also display the privileges associated with each role named in the clause:

```
mysql> SHOW GRANTS FOR 'u1'@'localhost' USING 'r1';
+-----------------------------------------+
| Grants for u1@localhost                 |
+-----------------------------------------+
| GRANT USAGE ON *.* TO `u1`@`localhost`  |
| GRANT SELECT ON `db1`.* TO `u1`@`localhost` |
| GRANT `r1`@`%`,`r2`@`%` TO `u1`@`localhost` |
+-----------------------------------------+
mysql> SHOW GRANTS FOR 'u1'@'localhost' USING 'r2';
+-----------------------------------------------------------+
| Grants for u1@localhost                                   |
+-----------------------------------------------------------+
| GRANT USAGE ON *.* TO `u1`@`localhost`                    |
| GRANT INSERT, UPDATE, DELETE ON `db1`.* TO `u1`@`localhost` |
| GRANT `r1`@`%`,`r2`@`%` TO `u1`@`localhost`               |
+-----------------------------------------------------------+
mysql> SHOW GRANTS FOR 'u1'@'localhost' USING 'r1', 'r2';
```

```
+-----------------------------------------------------------------------+
| Grants for u1@localhost                                               |
+-----------------------------------------------------------------------+
| GRANT USAGE ON *.* TO `u1`@`localhost`                                |
| GRANT SELECT, INSERT, UPDATE, DELETE ON `db1`.* TO `u1`@`localhost` |
| GRANT `r1`@`%`,`r2`@`%` TO `u1`@`localhost`                           |
+-----------------------------------------------------------------------+
```

Note

Just because a role is granted to an account does not mean the role is active. The active roles for an account can differ across and within sessions, depending on the account default roles and whether `SET ROLE` has been executed within a session.

`SHOW GRANTS` displays only the privileges granted explicitly to the named account. Other privileges that might be available to the account are not displayed. For example, if an anonymous account exists, the named account might be able to use its privileges, but `SHOW GRANTS` will not display them.

To display nonprivilege information for MySQL accounts, use the `SHOW CREATE USER` statement. See Section 13.7.5.12, “SHOW CREATE USER Syntax”.

13.7.5.22 SHOW INDEX Syntax

```
SHOW [EXTENDED] {INDEX | INDEXES | KEYS}
    {FROM | IN} tbl_name
    [{FROM | IN} db_name]
    [WHERE expr]
```

`SHOW INDEX` returns table index information. The format resembles that of the `SQLStatistics` call in ODBC. This statement requires some privilege for any column in the table.

You can use *`db_name.tbl_name`* as an alternative to the *`tbl_name FROM db_name`* syntax. These two statements are equivalent:

```
SHOW INDEX FROM mytable FROM mydb;
SHOW INDEX FROM mydb.mytable;
```

As of MySQL 8.0.2, the optional `EXTENDED` keyword causes the output to include information about hidden indexes that MySQL uses internally and are not accessible by users.

The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

```
mysql> SHOW INDEX FROM City\G
*************************** 1. row ***************************
        Table: city
   Non_unique: 0
     Key_name: PRIMARY
 Seq_in_index: 1
  Column_name: ID
    Collation: A
  Cardinality: 4321
     Sub_part: NULL
       Packed: NULL
         Null:
   Index_type: BTREE
      Comment:
Index_comment:
*************************** 2. row ***************************
        Table: city
   Non_unique: 1
     Key_name: CountryCode
 Seq_in_index: 1
```

```
  Column_name: CountryCode
    Collation: A
  Cardinality: 4321
     Sub_part: NULL
       Packed: NULL
         Null:
   Index_type: BTREE
      Comment:
Index_comment:
```

`SHOW INDEX` returns the following fields:

- `Table`

  The name of the table.

- `Non_unique`

  0 if the index cannot contain duplicates, 1 if it can.

- `Key_name`

  The name of the index. If the index is the primary key, the name is always `PRIMARY`.

- `Seq_in_index`

  The column sequence number in the index, starting with 1.

- `Column_name`

  The column name.

- `Collation`

  How the column is sorted in the index. This can have values `A` (ascending), `D` (descending), or `NULL` (not sorted).

- `Cardinality`

  An estimate of the number of unique values in the index. To update this number, run `ANALYZE TABLE` or (for `MyISAM` tables) `myisamchk -a`.

  `Cardinality` is counted based on statistics stored as integers, so the value is not necessarily exact even for small tables. The higher the cardinality, the greater the chance that MySQL uses the index when doing joins.

- `Sub_part`

  The index prefix. That is, the number of indexed characters if the column is only partly indexed, `NULL` if the entire column is indexed.

  Note

  Prefix limits are measured in bytes, whereas the prefix length in `CREATE TABLE`, `ALTER TABLE`, and `CREATE INDEX` statements is interpreted as number of characters for nonbinary string types (`CHAR`, `VARCHAR`, `TEXT`) and number of bytes for binary string types (`BINARY`, `VARBINARY`, `BLOB`). Take this into account when specifying a prefix length for a nonbinary string column that uses a multibyte character set.

  For additional information about index prefixes, see Section 8.3.4, “Column Indexes”, and Section 13.1.12, “CREATE INDEX Syntax”.

- `Packed`

Indicates how the key is packed. `NULL` if it is not.

- `Null`

  Contains `YES` if the column may contain `NULL` values and `''` if not.

- `Index_type`

  The index method used (`BTREE`, `FULLTEXT`, `HASH`, `RTREE`).

- `Comment`

  Information about the index not described in its own column, such as `disabled` if the index is disabled.

- `Index_comment`

  Any comment provided for the index with a `COMMENT` attribute when the index was created.

- `Visible`

  Whether the index is visible to the optimizer. See Section 8.3.11, “Invisible Indexes”.

You can also obtain information about table indexes from `INFORMATION_SCHEMA`, which contains a `STATISTICS` table. See Section 24.21, “The INFORMATION_SCHEMA STATISTICS Table”. The extended information about hidden indexes is available only using `SHOW EXTENDED INDEX`; it cannot be obtained from the `STATISTICS` table.

You can list a table's indexes with the `mysqlshow -k db_name tbl_name` command.

13.7.5.23 SHOW MASTER STATUS Syntax

```
SHOW MASTER STATUS
```

This statement provides status information about the binary log files of the master. It requires either the `SUPER` or `REPLICATION CLIENT` privilege.

Example:

```
mysql> SHOW MASTER STATUS\G
*************************** 1. row ***************************
             File: master-bin.000002
         Position: 1307
     Binlog_Do_DB: test
 Binlog_Ignore_DB: manual, mysql
Executed_Gtid_Set: 3E11FA47-71CA-11E1-9E33-C80AA9429562:1-5
1 row in set (0.00 sec)
```

When global transaction IDs are in use, `Executed_Gtid_Set` shows the set of GTIDs for transactions that have been executed on the master. This is the same as the value for the `gtid_executed` system variable on this server, as well as the value for `Executed_Gtid_Set` in the output of `SHOW SLAVE STATUS` on this server.

13.7.5.24 SHOW OPEN TABLES Syntax

```
SHOW OPEN TABLES
    [{FROM | IN} db_name]
    [LIKE 'pattern' | WHERE expr]
```

`SHOW OPEN TABLES` lists the non-`TEMPORARY` tables that are currently open in the table cache. See Section 8.4.3.1, “How MySQL Opens and Closes Tables”. The `FROM` clause, if present, restricts the

tables shown to those present in the *db_name* database. The `LIKE` clause, if present, indicates which table names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

`SHOW OPEN TABLES` output has the following columns:

- `Database`

  The database containing the table.

- `Table`

  The table name.

- `In_use`

  The number of table locks or lock requests there are for the table. For example, if one client acquires a lock for a table using `LOCK TABLE t1 WRITE`, `In_use` will be 1. If another client issues `LOCK TABLE t1 WRITE` while the table remains locked, the client will block waiting for the lock, but the lock request causes `In_use` to be 2. If the count is zero, the table is open but not currently being used. `In_use` is also increased by the `HANDLER ... OPEN` statement and decreased by `HANDLER ... CLOSE`.

- `Name_locked`

  Whether the table name is locked. Name locking is used for operations such as dropping or renaming tables.

If you have no privileges for a table, it does not show up in the output from `SHOW OPEN TABLES`.

13.7.5.25 SHOW PLUGINS Syntax

```
SHOW PLUGINS
```

`SHOW PLUGINS` displays information about server plugins. Plugin information is also available in the `INFORMATION_SCHEMA.PLUGINS` table. See Section 24.14, “The INFORMATION_SCHEMA PLUGINS Table”.

Example of `SHOW PLUGINS` output:

```
mysql> SHOW PLUGINS\G
*************************** 1. row ***************************
   Name: binlog
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: NULL
License: GPL
*************************** 2. row ***************************
   Name: CSV
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: NULL
License: GPL
*************************** 3. row ***************************
   Name: MEMORY
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: NULL
License: GPL
*************************** 4. row ***************************
   Name: MyISAM
 Status: ACTIVE
   Type: STORAGE ENGINE
Library: NULL
License: GPL
```

```
...
```

`SHOW PLUGINS` output has the following columns:

- `Name`: The name used to refer to the plugin in statements such as `INSTALL PLUGIN` and `UNINSTALL PLUGIN`.
- `Status`: The plugin status, one of `ACTIVE`, `INACTIVE`, `DISABLED`, or `DELETED`.
- `Type`: The type of plugin, such as `STORAGE ENGINE`, `INFORMATION_SCHEMA`, or `AUTHENTICATION`.
- `Library`: The name of the plugin shared library file. This is the name used to refer to the plugin file in statements such as `INSTALL PLUGIN` and `UNINSTALL PLUGIN`. This file is located in the directory named by the `plugin_dir` system variable. If the library name is `NULL`, the plugin is compiled in and cannot be uninstalled with `UNINSTALL PLUGIN`.
- `License`: How the plugin is licensed; for example, `GPL`.

For plugins installed with `INSTALL PLUGIN`, the `Name` and `Library` values are also registered in the `mysql.plugin` table.

For information about plugin data structures that form the basis of the information displayed by `SHOW PLUGINS`, see Section 28.2, “The MySQL Plugin API”.

13.7.5.26 SHOW PRIVILEGES Syntax

```
SHOW PRIVILEGES
```

`SHOW PRIVILEGES` shows the list of system privileges that the MySQL server supports. The privileges displayed include all static privileges, and all currently registered dynamic privileges.

```
mysql> SHOW PRIVILEGES\G
*************************** 1. row ***************************
Privilege: Alter
  Context: Tables
  Comment: To alter the table
*************************** 2. row ***************************
Privilege: Alter routine
  Context: Functions,Procedures
  Comment: To alter or drop stored functions/procedures
*************************** 3. row ***************************
Privilege: Create
  Context: Databases,Tables,Indexes
  Comment: To create new databases and tables
*************************** 4. row ***************************
Privilege: Create routine
  Context: Databases
  Comment: To use CREATE FUNCTION/PROCEDURE
*************************** 5. row ***************************
Privilege: Create temporary tables
  Context: Databases
  Comment: To use CREATE TEMPORARY TABLE
...
```

Privileges belonging to a specific user are displayed by the `SHOW GRANTS` statement. See Section 13.7.5.21, “SHOW GRANTS Syntax”, for more information.

13.7.5.27 SHOW PROCEDURE CODE Syntax

```
SHOW PROCEDURE CODE proc_name
```

This statement is a MySQL extension that is available only for servers that have been built with debugging support. It displays a representation of the internal implementation of the named stored

procedure. A similar statement, `SHOW FUNCTION CODE`, displays information about stored functions (see Section 13.7.5.19, “SHOW FUNCTION CODE Syntax”).

To use either statement, you must have the global `SELECT` privilege.

If the named routine is available, each statement produces a result set. Each row in the result set corresponds to one “instruction” in the routine. The first column is `Pos`, which is an ordinal number beginning with 0. The second column is `Instruction`, which contains an SQL statement (usually changed from the original source), or a directive which has meaning only to the stored-routine handler.

```
mysql> DELIMITER //
mysql> CREATE PROCEDURE p1 ()
    -> BEGIN
    ->   DECLARE fanta INT DEFAULT 55;
    ->   DROP TABLE t2;
    ->   LOOP
    ->     INSERT INTO t3 VALUES (fanta);
    ->     END LOOP;
    ->   END//
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW PROCEDURE CODE p1//
+-----+----------------------------------------+
| Pos | Instruction                            |
+-----+----------------------------------------+
|   0 | set fanta@0 55                         |
|   1 | stmt 9 "DROP TABLE t2"                 |
|   2 | stmt 5 "INSERT INTO t3 VALUES (fanta)" |
|   3 | jump 2                                 |
+-----+----------------------------------------+
4 rows in set (0.00 sec)
```

In this example, the nonexecutable `BEGIN` and `END` statements have disappeared, and for the `DECLARE` *`variable_name`* statement, only the executable part appears (the part where the default is assigned). For each statement that is taken from source, there is a code word `stmt` followed by a type (9 means `DROP`, 5 means `INSERT`, and so on). The final row contains an instruction `jump 2`, meaning `GOTO instruction #2`.

13.7.5.28 SHOW PROCEDURE STATUS Syntax

```
SHOW PROCEDURE STATUS
    [LIKE 'pattern' | WHERE expr]
```

This statement is a MySQL extension. It returns characteristics of a stored procedure, such as the database, name, type, creator, creation and modification dates, and character set information. A similar statement, `SHOW FUNCTION STATUS`, displays information about stored functions (see Section 13.7.5.20, “SHOW FUNCTION STATUS Syntax”).

The `LIKE` clause, if present, indicates which procedure or function names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

```
mysql> SHOW PROCEDURE STATUS LIKE 'sp1'\G
*************************** 1. row ***************************
                  Db: test
                Name: sp1
                Type: PROCEDURE
             Definer: testuser@localhost
            Modified: 2004-08-03 15:29:37
             Created: 2004-08-03 15:29:37
       Security_type: DEFINER
             Comment:
character_set_client: utf8
collation_connection: utf8_general_ci
```

```
  Database Collation: utf8mb4_0900_ai_ci
```

`character_set_client` is the session value of the `character_set_client` system variable when the routine was created. `collation_connection` is the session value of the `collation_connection` system variable when the routine was created. `Database Collation` is the collation of the database with which the routine is associated.

You can also get information about stored routines from the `ROUTINES` table in `INFORMATION_SCHEMA`. See Section 24.18, “The INFORMATION_SCHEMA ROUTINES Table”.

13.7.5.29 SHOW PROCESSLIST Syntax

```
SHOW [FULL] PROCESSLIST
```

`SHOW PROCESSLIST` shows you which threads are running. You can also get this information from the `INFORMATION_SCHEMA PROCESSLIST` table or the `mysqladmin processlist` command. If you have the `PROCESS` privilege, you can see all threads. Otherwise, you can see only your own threads (that is, threads associated with the MySQL account that you are using). If you do not use the `FULL` keyword, only the first 100 characters of each statement are shown in the `Info` field.

Process information is also available from the `performance_schema.threads` table. However, access to `threads` does not require a mutex and has minimal impact on server performance. `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` have negative performance consequences because they require a mutex. `threads` also shows information about background threads, which `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` do not. This means that `threads` can be used to monitor activity the other thread information sources cannot.

The `SHOW PROCESSLIST` statement is very useful if you get the “too many connections” error message and want to find out what is going on. MySQL reserves one extra connection to be used by accounts that have the `SUPER` privilege, to ensure that administrators should always be able to connect and check the system (assuming that you are not giving this privilege to all your users).

Threads can be killed with the `KILL` statement. See Section 13.7.6.4, “KILL Syntax”.

Here is an example of `SHOW PROCESSLIST` output:

```
mysql> SHOW FULL PROCESSLIST\G
*************************** 1. row ***************************
Id: 1
User: system user
Host:
db: NULL
Command: Connect
Time: 1030455
State: Waiting for master to send event
Info: NULL
*************************** 2. row ***************************
Id: 2
User: system user
Host:
db: NULL
Command: Connect
Time: 1004
State: Has read all relay log; waiting for the slave
       I/O thread to update it
Info: NULL
*************************** 3. row ***************************
Id: 3112
User: replikator
Host: artemis:2204
db: NULL
Command: Binlog Dump
Time: 2144
State: Has sent all binlog to slave; waiting for binlog to be updated
```

```
Info: NULL
*************************** 4. row ***************************
Id: 3113
User: replikator
Host: iconnect2:45781
db: NULL
Command: Binlog Dump
Time: 2086
State: Has sent all binlog to slave; waiting for binlog to be updated
Info: NULL
*************************** 5. row ***************************
Id: 3123
User: stefan
Host: localhost
db: apollon
Command: Query
Time: 0
State: NULL
Info: SHOW FULL PROCESSLIST
5 rows in set (0.00 sec)
```

The columns produced by `SHOW PROCESSLIST` have the following meanings:

- `Id`

  The connection identifier. This is the same type of value displayed in the `ID` column of the `INFORMATION_SCHEMA.PROCESSLIST` table, the `PROCESSLIST_ID` column of the Performance Schema `threads` table, and returned by the `CONNECTION_ID()` function.

- `User`

  The MySQL user who issued the statement. If this is `system user`, it refers to a nonclient thread spawned by the server to handle tasks internally. This could be the I/O or SQL thread used on replication slaves or a delayed-row handler. `unauthenticated user` refers to a thread that has become associated with a client connection but for which authentication of the client user has not yet been done. `event_scheduler` refers to the thread that monitors scheduled events. For `system user`, there is no host specified in the `Host` column.

- `Host`

  The host name of the client issuing the statement (except for `system user` where there is no host). `SHOW PROCESSLIST` reports the host name for TCP/IP connections in *`host_name:client_port`* format to make it easier to determine which client is doing what.

- `db`

  The default database, if one is selected, otherwise `NULL`.

- `Command`

  The type of command the thread is executing. For descriptions for thread commands, see Section 8.14, “Examining Thread Information”. The value of this column corresponds to the `COM_`*`xxx`* commands of the client/server protocol and `Com_`*`xxx`* status variables. See Section 5.1.7, “Server Status Variables”

- `Time`

  The time in seconds that the thread has been in its current state. For a slave SQL thread, the value is the number of seconds between the timestamp of the last replicated event and the real time of the slave machine. See Section 18.2.2, “Replication Implementation Details”.

- `State`

  An action, event, or state that indicates what the thread is doing. Descriptions for `State` values can be found at Section 8.14, “Examining Thread Information”.

Most states correspond to very quick operations. If a thread stays in a given state for many seconds, there might be a problem that needs to be investigated.

For the `SHOW PROCESSLIST` statement, the value of `State` is `NULL`.

- `Info`

  The statement the thread is executing, or `NULL` if it is not executing any statement. The statement might be the one sent to the server, or an innermost statement if the statement executes other statements. For example, if a `CALL` statement executes a stored procedure that is executing a `SELECT` statement, the `Info` value shows the `SELECT` statement.

13.7.5.30 SHOW PROFILE Syntax

```
SHOW PROFILE [type [, type] ... ]
    [FOR QUERY n]
    [LIMIT row_count [OFFSET offset]]

type:
    ALL
  | BLOCK IO
  | CONTEXT SWITCHES
  | CPU
  | IPC
  | MEMORY
  | PAGE FAULTS
  | SOURCE
  | SWAPS
```

The `SHOW PROFILE` and `SHOW PROFILES` statements display profiling information that indicates resource usage for statements executed during the course of the current session.

Note

These statements are deprecated and will be removed in a future MySQL release. Use the Performance Schema instead; see Section 25.18.1, “Query Profiling Using Performance Schema”.

Profiling is controlled by the `profiling` session variable, which has a default value of 0 (`OFF`). Profiling is enabled by setting `profiling` to 1 or `ON`:

```
mysql> SET profiling = 1;
```

`SHOW PROFILES` displays a list of the most recent statements sent to the server. The size of the list is controlled by the `profiling_history_size` session variable, which has a default value of 15. The maximum value is 100. Setting the value to 0 has the practical effect of disabling profiling.

All statements are profiled except `SHOW PROFILE` and `SHOW PROFILES`, so you will find neither of those statements in the profile list. Malformed statements are profiled. For example, `SHOW PROFILING` is an illegal statement, and a syntax error occurs if you try to execute it, but it will show up in the profiling list.

`SHOW PROFILE` displays detailed information about a single statement. Without the `FOR QUERY n` clause, the output pertains to the most recently executed statement. If `FOR QUERY n` is included, `SHOW PROFILE` displays information for statement *`n`*. The values of *`n`* correspond to the `Query_ID` values displayed by `SHOW PROFILES`.

The `LIMIT row_count` clause may be given to limit the output to *`row_count`* rows. If `LIMIT` is given, `OFFSET offset` may be added to begin the output *`offset`* rows into the full set of rows.

By default, `SHOW PROFILE` displays `Status` and `Duration` columns. The `Status` values are like the `State` values displayed by `SHOW PROCESSLIST`, although there might be some minor differences

in interpretion for the two statements for some status values (see Section 8.14, “Examining Thread Information”).

Optional *`type`* values may be specified to display specific additional types of information:

- `ALL` displays all information
- `BLOCK IO` displays counts for block input and output operations
- `CONTEXT SWITCHES` displays counts for voluntary and involuntary context switches
- `CPU` displays user and system CPU usage times
- `IPC` displays counts for messages sent and received
- `MEMORY` is not currently implemented
- `PAGE FAULTS` displays counts for major and minor page faults
- `SOURCE` displays the names of functions from the source code, together with the name and line number of the file in which the function occurs
- `SWAPS` displays swap counts

Profiling is enabled per session. When a session ends, its profiling information is lost.

```
mysql> SELECT @@profiling;
+-------------+
| @@profiling |
+-------------+
|           0 |
+-------------+
1 row in set (0.00 sec)

mysql> SET profiling = 1;
Query OK, 0 rows affected (0.00 sec)

mysql> DROP TABLE IF EXISTS t1;
Query OK, 0 rows affected, 1 warning (0.00 sec)

mysql> CREATE TABLE T1 (id INT);
Query OK, 0 rows affected (0.01 sec)

mysql> SHOW PROFILES;
+----------+----------+--------------------------+
| Query_ID | Duration | Query                    |
+----------+----------+--------------------------+
|        0 | 0.000088 | SET PROFILING = 1        |
|        1 | 0.000136 | DROP TABLE IF EXISTS t1  |
|        2 | 0.011947 | CREATE TABLE t1 (id INT) |
+----------+----------+--------------------------+
3 rows in set (0.00 sec)

mysql> SHOW PROFILE;
+----------------------+----------+
| Status               | Duration |
+----------------------+----------+
| checking permissions | 0.000040 |
| creating table       | 0.000056 |
| After create         | 0.011363 |
| query end            | 0.000375 |
| freeing items        | 0.000089 |
| logging slow query   | 0.000019 |
| cleaning up          | 0.000005 |
+----------------------+----------+
7 rows in set (0.00 sec)

mysql> SHOW PROFILE FOR QUERY 1;
+--------------------+----------+
```

```
| Status             | Duration |
+--------------------+----------+
| query end          | 0.000107 |
| freeing items      | 0.000008 |
| logging slow query | 0.000015 |
| cleaning up        | 0.000006 |
+--------------------+----------+
4 rows in set (0.00 sec)

mysql> SHOW PROFILE CPU FOR QUERY 2;
+----------------------+----------+----------+------------+
| Status               | Duration | CPU_user | CPU_system |
+----------------------+----------+----------+------------+
| checking permissions | 0.000040 | 0.000038 |   0.000002 |
| creating table       | 0.000056 | 0.000028 |   0.000028 |
| After create         | 0.011363 | 0.000217 |   0.001571 |
| query end            | 0.000375 | 0.000013 |   0.000028 |
| freeing items        | 0.000089 | 0.000010 |   0.000014 |
| logging slow query   | 0.000019 | 0.000009 |   0.000010 |
| cleaning up          | 0.000005 | 0.000003 |   0.000002 |
+----------------------+----------+----------+------------+
7 rows in set (0.00 sec)
```

Note

Profiling is only partially functional on some architectures. For values that depend on the `getrusage()` system call, `NULL` is returned on systems such as Windows that do not support the call. In addition, profiling is per process and not per thread. This means that activity on threads within the server other than your own may affect the timing information that you see.

You can also get profiling information from the `PROFILING` table in `INFORMATION_SCHEMA`. See Section 24.16, “The INFORMATION_SCHEMA PROFILING Table”. For example, the following queries produce the same result:

```
SHOW PROFILE FOR QUERY 2;

SELECT STATE, FORMAT(DURATION, 6) AS DURATION
FROM INFORMATION_SCHEMA.PROFILING
WHERE QUERY_ID = 2 ORDER BY SEQ;
```

13.7.5.31 SHOW PROFILES Syntax

```
SHOW PROFILES
```

The `SHOW PROFILES` statement, together with `SHOW PROFILE`, displays profiling information that indicates resource usage for statements executed during the course of the current session. For more information, see Section 13.7.5.30, “SHOW PROFILE Syntax”.

Note

These statements are deprecated and will be removed in a future MySQL release. Use the Performance Schema instead; see Chapter 25, *MySQL Performance Schema*.

13.7.5.32 SHOW RELAYLOG EVENTS Syntax

```
SHOW RELAYLOG EVENTS
    [IN 'log_name']
    [FROM pos]
    [LIMIT [offset,] row_count]
    [channel_option]

channel_option:
    FOR CHANNEL channel
```

Shows the events in the relay log of a replication slave. If you do not specify `'log_name'`, the first relay log is displayed. This statement has no effect on the master.

The `LIMIT` clause has the same syntax as for the `SELECT` statement. See Section 13.2.10, "SELECT Syntax".

Note

Issuing a `SHOW RELAYLOG EVENTS` with no `LIMIT` clause could start a very time- and resource-consuming process because the server returns to the client the complete contents of the relay log (including all statements modifying data that have been received by the slave).

The optional `FOR CHANNEL channel` clause enables you to choose which replication channel the statement applies to. If no clause is set and no extra channels exist, the statement applies to the default channel. Providing a `FOR CHANNEL channel` clause applies the `SHOW RELAYLOG EVENTS` statement to a specific replication channel.

When using multiple replication channels, if a `SHOW RELAYLOG EVENTS` statement does not have a channel defined using a `FOR CHANNEL channel` clause an error is generated. See Section 18.2.3, "Replication Channels" for more information.

Note

Some events relating to the setting of user and system variables are not included in the output from `SHOW RELAYLOG EVENTS`. To get complete coverage of events within a relay log, use `mysqlbinlog`.

13.7.5.33 SHOW SLAVE HOSTS Syntax

```
SHOW SLAVE HOSTS
```

Displays a list of replication slaves currently registered with the master.

`SHOW SLAVE HOSTS` should be executed on a server that acts as a replication master. The statement displays information about servers that are or have been connected as replication slaves, with each row of the result corresponding to one slave server, as shown here:

```
mysql> SHOW SLAVE HOSTS;
+------------+-----------+------+-----------+--------------------------------------+
| Server_id  | Host      | Port | Master_id | Slave_UUID                           |
+------------+-----------+------+-----------+--------------------------------------+
|  192168010 | iconnect2 | 3306 | 192168011 | 14cb6624-7f93-11e0-b2c0-c80aa9429562 |
| 1921680101 | athena    | 3306 | 192168011 | 07af4990-f41f-11df-a566-7ac56fdaf645 |
+------------+-----------+------+-----------+--------------------------------------+
```

- `Server_id`: The unique server ID of the slave server, as configured in the slave server's option file, or on the command line with `--server-id=value`.
- `Host`: The host name of the slave server as specified on the slave with the `--report-host` option. This can differ from the machine name as configured in the operating system.
- `User`: The slave server user name as, specified on the slave with the `--report-user` option. Statement output includes this column only if the master server is started with the `--show-slave-auth-info` option.
- `Password`: The slave server password as, specified on the slave with the `--report-password` option. Statement output includes this column only if the master server is started with the `--show-slave-auth-info` option.
- `Port`: The port on the master to which the slave server is listening, as specified on the slave with the `--report-port` option.

A zero in this column means that the slave port (`--report-port`) was not set.

- `Master_id`: The unique server ID of the master server that the slave server is replicating from. This is the server ID of the server on which `SHOW SLAVE HOSTS` is executed, so this same value is listed for each row in the result.
- `Slave_UUID`: The globally unique ID of this slave, as generated on the slave and found in the slave's `auto.cnf` file.

13.7.5.34 SHOW SLAVE STATUS Syntax

```
SHOW SLAVE STATUS [FOR CHANNEL channel]
```

This statement provides status information on essential parameters of the slave threads. It requires either the `SUPER` or `REPLICATION CLIENT` privilege.

`SHOW SLAVE STATUS` is nonblocking. When run concurrently with `STOP SLAVE`, `SHOW SLAVE STATUS` returns without waiting for `STOP SLAVE` to finish shutting down the slave SQL thread or slave I/O thread (or both). This permits use in monitoring and other applications where getting an immediate response from `SHOW SLAVE STATUS` more important than ensuring that it returned the latest data.

If you issue this statement using the `mysql` client, you can use a `\G` statement terminator rather than a semicolon to obtain a more readable vertical layout:

```
mysql> SHOW SLAVE STATUS\G
*************************** 1. row ***************************
               Slave_IO_State: Waiting for master to send event
                  Master_Host: localhost
                  Master_User: root
                  Master_Port: 13000
                Connect_Retry: 60
              Master_Log_File: master-bin.000002
          Read_Master_Log_Pos: 1307
               Relay_Log_File: slave-relay-bin.000003
                Relay_Log_Pos: 1508
        Relay_Master_Log_File: master-bin.000002
             Slave_IO_Running: Yes
            Slave_SQL_Running: Yes
              Replicate_Do_DB:
          Replicate_Ignore_DB:
           Replicate_Do_Table:
       Replicate_Ignore_Table:
      Replicate_Wild_Do_Table:
  Replicate_Wild_Ignore_Table:
                   Last_Errno: 0
                   Last_Error:
                 Skip_Counter: 0
          Exec_Master_Log_Pos: 1307
              Relay_Log_Space: 1858
              Until_Condition: None
               Until_Log_File:
                Until_Log_Pos: 0
           Master_SSL_Allowed: No
           Master_SSL_CA_File:
           Master_SSL_CA_Path:
              Master_SSL_Cert:
            Master_SSL_Cipher:
               Master_SSL_Key:
        Seconds_Behind_Master: 0
Master_SSL_Verify_Server_Cert: No
                Last_IO_Errno: 0
                Last_IO_Error:
               Last_SQL_Errno: 0
               Last_SQL_Error:
  Replicate_Ignore_Server_Ids:
```

```
             Master_Server_Id: 1
                  Master_UUID: 3e11fa47-71ca-11e1-9e33-c80aa9429562
             Master_Info_File: /var/mysqld.2/data/master.info
                    SQL_Delay: 0
          SQL_Remaining_Delay: NULL
      Slave_SQL_Running_State: Slave has read all relay log; waiting for the slave I/O thread to update
           Master_Retry_Count: 10
                  Master_Bind:
      Last_IO_Error_Timestamp:
     Last_SQL_Error_Timestamp:
               Master_SSL_Crl:
           Master_SSL_Crlpath:
           Retrieved_Gtid_Set: 3e11fa47-71ca-11e1-9e33-c80aa9429562:1-5
            Executed_Gtid_Set: 3e11fa47-71ca-11e1-9e33-c80aa9429562:1-5
                Auto_Position: 1
         Replicate_Rewrite_DB:
                 Channel_name:
1 row in set (0.00 sec)
```

The Performance Schema provides tables that expose replication information. This is similar to the information available from the `SHOW SLAVE STATUS` statement, but represented in table form. For details, see Section 25.11.11, “Performance Schema Replication Tables”.

The following list describes the fields returned by `SHOW SLAVE STATUS`. For additional information about interpreting their meanings, see Section 18.1.7.1, “Checking Replication Status”.

- `Slave_IO_State`

  A copy of the `State` field of the `SHOW PROCESSLIST` output for the slave I/O thread. This tells you what the thread is doing: trying to connect to the master, waiting for events from the master, reconnecting to the master, and so on. For a listing of possible states, see Section 8.14.5, “Replication Slave I/O Thread States”.

- `Master_Host`

  The master host that the slave is connected to.

- `Master_User`

  The user name of the account used to connect to the master.

- `Master_Port`

  The port used to connect to the master.

- `Connect_Retry`

  The number of seconds between connect retries (default 60). This can be set with the `CHANGE MASTER TO` statement.

- `Master_Log_File`

  The name of the master binary log file from which the I/O thread is currently reading.

- `Read_Master_Log_Pos`

  The position in the current master binary log file up to which the I/O thread has read.

- `Relay_Log_File`

  The name of the relay log file from which the SQL thread is currently reading and executing.

- `Relay_Log_Pos`

  The position in the current relay log file up to which the SQL thread has read and executed.

- `Relay_Master_Log_File`

  The name of the master binary log file containing the most recent event executed by the SQL thread.

- `Slave_IO_Running`

  Whether the I/O thread is started and has connected successfully to the master. Internally, the state of this thread is represented by one of the following three values:

  - **MYSQL_SLAVE_NOT_RUN.** The slave I/O thread is not running. For this state, `Slave_IO_Running` is `No`.
  - **MYSQL_SLAVE_RUN_NOT_CONNECT.** The slave I/O thread is running, but is not connected to a replication master. For this state, `Slave_IO_Running` depends on the server version as shown in the following table.

MySQL Version	`Slave_IO_Running`
4.1 (4.1.13 and earlier); 5.0 (5.0.11 and earlier)	`Yes`
4.1 (4.1.14 and later); 5.0 (5.0.12 and later)	`No`
5.1 (5.1.45 and earlier)	`No`
5.1 (5.1.46 and later); 5.5; 5.6	`Connecting`

  - **MYSQL_SLAVE_RUN_CONNECT.** The slave I/O thread is running, and is connected to a replication master. For this state, `Slave_IO_Running` is `Yes`.

  The value of the `Slave_running` system status variable corresponds with this value.

- `Slave_SQL_Running`

  Whether the SQL thread is started.

- `Replicate_Do_DB`, `Replicate_Ignore_DB`

  The lists of databases that were specified with the `--replicate-do-db` and `--replicate-ignore-db` options, if any.

- `Replicate_Do_Table`, `Replicate_Ignore_Table`, `Replicate_Wild_Do_Table`, `Replicate_Wild_Ignore_Table`

  The lists of tables that were specified with the `--replicate-do-table`, `--replicate-ignore-table`, `--replicate-wild-do-table`, and `--replicate-wild-ignore-table` options, if any.

- `Last_Errno`, `Last_Error`

  These columns are aliases for `Last_SQL_Errno` and `Last_SQL_Error`.

  Issuing `RESET MASTER` or `RESET SLAVE` resets the values shown in these columns.

  Note

  When the slave SQL thread receives an error, it reports the error first, then stops the SQL thread. This means that there is a small window of time during which `SHOW SLAVE STATUS` shows a nonzero value for `Last_SQL_Errno` even though `Slave_SQL_Running` still displays `Yes`.

- `Skip_Counter`

  The current value of the `sql_slave_skip_counter` system variable. See Section 13.4.2.5, “SET GLOBAL sql_slave_skip_counter Syntax”.

- `Exec_Master_Log_Pos`

  The position in the current master binary log file to which the SQL thread has read and executed, marking the start of the next transaction or event to be processed. You can use this value with the `CHANGE MASTER TO` statement's `MASTER_LOG_POS` option when starting a new slave from an existing slave, so that the new slave reads from this point. The coordinates given by (`Relay_Master_Log_File`, `Exec_Master_Log_Pos`) in the master's binary log correspond to the coordinates given by (`Relay_Log_File`, `Relay_Log_Pos`) in the relay log.

  Inconsistencies in the sequence of transactions from the relay log which have been executed can cause this value to be a “low-water mark”. In other words, transactions appearing before the position are guaranteed to have committed, but transactions after the position may have committed or not. If these gaps need to be corrected, use `START SLAVE UNTIL SQL_AFTER_MTS_GAPS`. See Section 18.4.1.34, “Replication and Transaction Inconsistencies” for more information.

- `Relay_Log_Space`

  The total combined size of all existing relay log files.

- `Until_Condition`, `Until_Log_File`, `Until_Log_Pos`

  The values specified in the `UNTIL` clause of the `START SLAVE` statement.

  `Until_Condition` has these values:

  - `None` if no `UNTIL` clause was specified
  - `Master` if the slave is reading until a given position in the master's binary log
  - `Relay` if the slave is reading until a given position in its relay log
  - `SQL_BEFORE_GTIDS` if the slave SQL thread is processing transactions until it has reached the first transaction whose GTID is listed in the `gtid_set`.
  - `SQL_AFTER_GTIDS` if the slave threads are processing all transactions until the last transaction in the `gtid_set` has been processed by both threads.
  - `SQL_AFTER_MTS_GAPS` if a multi-threaded slave's SQL threads are running until no more gaps are found in the relay log.

  `Until_Log_File` and `Until_Log_Pos` indicate the log file name and position that define the coordinates at which the SQL thread stops executing.

  For more information on `UNTIL` clauses, see Section 13.4.2.6, “START SLAVE Syntax”.

- `Master_SSL_Allowed`, `Master_SSL_CA_File`, `Master_SSL_CA_Path`, `Master_SSL_Cert`, `Master_SSL_Cipher`, `Master_SSL_CRL_File`, `Master_SSL_CRL_Path`, `Master_SSL_Key`, `Master_SSL_Verify_Server_Cert`

  These fields show the SSL parameters used by the slave to connect to the master, if any.

  `Master_SSL_Allowed` has these values:

  - `Yes` if an SSL connection to the master is permitted
  - `No` if an SSL connection to the master is not permitted
  - `Ignored` if an SSL connection is permitted but the slave server does not have SSL support enabled

  The values of the other SSL-related fields correspond to the values of the `MASTER_SSL_CA`, `MASTER_SSL_CAPATH`, `MASTER_SSL_CERT`, `MASTER_SSL_CIPHER`, `MASTER_SSL_CRL`,

MASTER_SSL_CRLPATH, MASTER_SSL_KEY, and MASTER_SSL_VERIFY_SERVER_CERT options to the CHANGE MASTER TO statement. See Section 13.4.2.1, “CHANGE MASTER TO Syntax”.

- Seconds_Behind_Master

  This field is an indication of how “late” the slave is:

  - When the slave is actively processing updates, this field shows the difference between the current timestamp on the slave and the original timestamp logged on the master for the event currently being processed on the slave.
  - When no event is currently being processed on the slave, this value is 0.

  In essence, this field measures the time difference in seconds between the slave SQL thread and the slave I/O thread. If the network connection between master and slave is fast, the slave I/O thread is very close to the master, so this field is a good approximation of how late the slave SQL thread is compared to the master. If the network is slow, this is *not* a good approximation; the slave SQL thread may quite often be caught up with the slow-reading slave I/O thread, so Seconds_Behind_Master often shows a value of 0, even if the I/O thread is late compared to the master. In other words, *this column is useful only for fast networks*.

  This time difference computation works even if the master and slave do not have identical clock times, provided that the difference, computed when the slave I/O thread starts, remains constant from then on. Any changes—including NTP updates—can lead to clock skews that can make calculation of Seconds_Behind_Master less reliable.

  In MySQL 8.0, this field is NULL (undefined or unknown) if the slave SQL thread is not running, or if the SQL thread has consumed all of the relay log and the slave I/O thread is not running. (In older versions of MySQL, this field was NULL if the slave SQL thread or the slave I/O thread was not running or was not connected to the master.) If the I/O thread is running but the relay log is exhausted, Seconds_Behind_Master is set to 0.

  The value of Seconds_Behind_Master is based on the timestamps stored in events, which are preserved through replication. This means that if a master M1 is itself a slave of M0, any event from M1's binary log that originates from M0's binary log has M0's timestamp for that event. This enables MySQL to replicate TIMESTAMP successfully. However, the problem for Seconds_Behind_Master is that if M1 also receives direct updates from clients, the Seconds_Behind_Master value randomly fluctuates because sometimes the last event from M1 originates from M0 and sometimes is the result of a direct update on M1.

  When using a multi-threaded slave, you should keep in mind that this value is based on Exec_Master_Log_Pos, and so may not reflect the position of the most recently committed transaction.

- Last_IO_Errno, Last_IO_Error

  The error number and error message of the most recent error that caused the I/O thread to stop. An error number of 0 and message of the empty string mean “no error.” If the Last_IO_Error value is not empty, the error values also appear in the slave's error log.

  I/O error information includes a timestamp showing when the most recent I/O thread error occurred. This timestamp uses the format YYMMDD HH:MM:SS, and appears in the Last_IO_Error_Timestamp column.

  Issuing RESET MASTER or RESET SLAVE resets the values shown in these columns.

- Last_SQL_Errno, Last_SQL_Error

  The error number and error message of the most recent error that caused the SQL thread to stop. An error number of 0 and message of the empty string mean “no error.” If the Last_SQL_Error value is not empty, the error values also appear in the slave's error log.

If the slave is multi-threaded, the SQL thread is the coordinator for worker threads. In this case, the `Last_SQL_Error` field shows exactly what the `Last_Error_Message` column in the Performance Schema `replication_applier_status_by_coordinator` table shows. The field value is modified to suggest that there may be more failures in the other worker threads which can be seen in the `replication_applier_status_by_worker` table that shows each worker thread's status. If that table is not available, the slave error log can be used. The log or the `replication_applier_status_by_worker` table should also be used to learn more about the failure shown by `SHOW SLAVE STATUS` or the coordinator table.

SQL error information includes a timestamp showing when the most recent SQL thread error occurred. This timestamp uses the format `YYMMDD HH:MM:SS`, and appears in the `Last_SQL_Error_Timestamp` column.

Issuing `RESET MASTER` or `RESET SLAVE` resets the values shown in these columns.

In MySQL 8.0, all error codes and messages displayed in the `Last_SQL_Errno` and `Last_SQL_Error` columns correspond to error values listed in Section B.3, “Server Error Codes and Messages”. This was not always true in previous versions. (Bug #11760365, Bug #52768)

- `Replicate_Ignore_Server_Ids`

  In MySQL 8.0, you set a slave to ignore events from 0 or more masters using the `IGNORE_SERVER_IDS` option of the `CHANGE MASTER TO` statement. By default this is blank, and is usually modified only when using a circular or other multi-master replication setup. The message shown for `Replicate_Ignore_Server_Ids` when not blank consists of a comma-delimited list of one or more numbers, indicating the server IDs to be ignored. For example:

  ```
  Replicate_Ignore_Server_Ids: 2, 6, 9
  ```

  Note

  `Ignored_server_ids` also shows the server IDs to be ignored, but is a space-delimited list, which is preceded by the total number of server IDs to be ignored. For example, if a `CHANGE MASTER TO` statement containing the `IGNORE_SERVER_IDS = (2,6,9)` option has been issued to tell a slave to ignore masters having the server ID 2, 6, or 9, that information appears as shown here:

  ```
  Ignored_server_ids: 3, 2, 6, 9
  ```

  The first number (in this case 3) shows the number of server IDs being ignored.

  `Replicate_Ignore_Server_Ids` filtering is performed by the I/O thread, rather than by the SQL thread, which means that events which are filtered out are not written to the relay log. This differs from the filtering actions taken by server options such `--replicate-do-table`, which apply to the SQL thread.

- `Master_Server_Id`

  The `server_id` value from the master.

- `Master_UUID`

  The `server_uuid` value from the master.

- `Master_Info_File`

  The location of the `master.info` file.

- `SQL_Delay`

  The number of seconds that the slave must lag the master.

- `SQL_Remaining_Delay`

  When `Slave_SQL_Running_State` is `Waiting until MASTER_DELAY seconds after master executed event`, this field contains the number of delay seconds remaining. At other times, this field is `NULL`.

- `Slave_SQL_Running_State`

  The state of the SQL thread (analogous to `Slave_IO_State`). The value is identical to the `State` value of the SQL thread as displayed by `SHOW PROCESSLIST`. Section 8.14.6, "Replication Slave SQL Thread States", provides a listing of possible states

- `Master_Retry_Count`

  The number of times the slave can attempt to reconnect to the master in the event of a lost connection. This value can be set using the `MASTER_RETRY_COUNT` option of the `CHANGE MASTER TO` statement (preferred) or the older `--master-retry-count` server option (still supported for backward compatibility).

- `Master_Bind`

  The network interface that the slave is bound to, if any. This is set using the `MASTER_BIND` option for the `CHANGE MASTER TO` statement.

- `Last_IO_Error_Timestamp`

  A timestamp in `YYMMDD HH:MM:SS` format that shows when the most recent I/O error took place.

- `Last_SQL_Error_Timestamp`

  A timestamp in `YYMMDD HH:MM:SS` format that shows when the last SQL error occurred.

- `Retrieved_Gtid_Set`

  The set of global transaction IDs corresponding to all transactions received by this slave. Empty if GTIDs are not in use. See GTID Sets for more information.

  This is the set of all GTIDs that exist or have existed in the relay logs. Each GTID is added as soon as the `Gtid_log_event` is received. This can cause partially transmitted transactions to have their GTIDs included in the set.

  When all relay logs are lost due to executing `RESET SLAVE` or `CHANGE MASTER TO`, or due to the effects of the `--relay-log-recovery` option, the set is cleared. When `relay_log_purge = 1`, the newest relay log is always kept, and the set is not cleared.

- `Executed_Gtid_Set`

  The set of global transaction IDs written in the binary log. This is the same as the value for the global `gtid_executed` system variable on this server, as well as the value for `Executed_Gtid_Set` in the output of `SHOW MASTER STATUS` on this server. Empty if GTIDs are not in use. See GTID Sets for more information.

- `Auto_Position`

  1 if autopositioning is in use; otherwise 0.

- `Replicate_Rewrite_DB`

The `Replicate_Rewrite_DB` value displays any replication filtering rules that were specified. For example, if the following replication filter rule was set:

```
CHANGE REPLICATION FILTER REPLICATE_REWRITE_DB=((db1,db2), (db3,db4));
```

the `Replicate_Rewrite_DB` value displays:

```
Replicate_Rewrite_DB: (db1,db2),(db3,db4)
```

For more information, see Section 13.4.2.2, “CHANGE REPLICATION FILTER Syntax”.

- `Channel_name`

  The replication channel which is being displayed. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

13.7.5.35 SHOW STATUS Syntax

```
SHOW [GLOBAL | SESSION] STATUS
    [LIKE 'pattern' | WHERE expr]
```

`SHOW STATUS` provides server status information (see Section 5.1.7, “Server Status Variables”). This statement does not require any privilege. It requires only the ability to connect to the server.

Status variable information is also available from these sources:

- Performance Schema tables. See Section 25.11.14, “Performance Schema Status Variable Tables”.
- The `mysqladmin extended-status` command. See Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”.

For `SHOW STATUS`, a `LIKE` clause, if present, indicates which variable names to match. A `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

`SHOW STATUS` accepts an optional `GLOBAL` or `SESSION` variable scope modifier:

- With a `GLOBAL` modifier, the statement displays the global status values. A global status variable may represent status for some aspect of the server itself (for example, `Aborted_connects`), or the aggregated status over all connections to MySQL (for example, `Bytes_received` and `Bytes_sent`). If a variable has no global value, the session value is displayed.
- With a `SESSION` modifier, the statement displays the status variable values for the current connection. If a variable has no session value, the global value is displayed. `LOCAL` is a synonym for `SESSION`.
- If no modifier is present, the default is `SESSION`.

The scope for each status variable is listed at Section 5.1.7, “Server Status Variables”.

Each invocation of the `SHOW STATUS` statement uses an internal temporary table and increments the global `Created_tmp_tables` value.

Partial output is shown here. The list of names and values may differ for your server. The meaning of each variable is given in Section 5.1.7, “Server Status Variables”.

```
mysql> SHOW STATUS;
+--------------------------+------------+
| Variable_name            | Value      |
```

```
+-------------------------+------------+
| Aborted_clients         | 0          |
| Aborted_connects        | 0          |
| Bytes_received          | 155372598  |
| Bytes_sent              | 1176560426 |
| Connections             | 30023      |
| Created_tmp_disk_tables | 0          |
| Created_tmp_tables      | 8340       |
| Created_tmp_files       | 60         |
...
| Open_tables             | 1          |
| Open_files              | 2          |
| Open_streams            | 0          |
| Opened_tables           | 44600      |
| Questions               | 2026873    |
...
| Table_locks_immediate   | 1920382    |
| Table_locks_waited      | 0          |
| Threads_cached          | 0          |
| Threads_created         | 30022      |
| Threads_connected       | 1          |
| Threads_running         | 1          |
| Uptime                  | 80380      |
+-------------------------+------------+
```

With a `LIKE` clause, the statement displays only rows for those variables with names that match the pattern:

```
mysql> SHOW STATUS LIKE 'Key%';
+--------------------+----------+
| Variable_name      | Value    |
+--------------------+----------+
| Key_blocks_used    | 14955    |
| Key_read_requests  | 96854827 |
| Key_reads          | 162040   |
| Key_write_requests | 7589728  |
| Key_writes         | 3813196  |
+--------------------+----------+
```

13.7.5.36 SHOW TABLE STATUS Syntax

```
SHOW TABLE STATUS
    [{FROM | IN} db_name]
    [LIKE 'pattern' | WHERE expr]
```

`SHOW TABLE STATUS` works likes `SHOW TABLES`, but provides a lot of information about each non-`TEMPORARY` table. You can also get this list using the `mysqlshow --status db_name` command. The `LIKE` clause, if present, indicates which table names to match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

This statement also displays information about views.

`SHOW TABLE STATUS` output has the following columns:

- `Name`

  The name of the table.

- `Engine`

  The storage engine for the table. See Chapter 16, *Alternative Storage Engines*.

- `Version`

  This field is unused.

- `Row_format`

  The row-storage format (`Fixed`, `Dynamic`, `Compressed`, `Redundant`, `Compact`). For `MyISAM` tables, (`Dynamic` corresponds to what `myisamchk -dvv` reports as `Packed`.

- `Rows`

  The number of rows. Some storage engines, such as `MyISAM`, store the exact count. For other storage engines, such as `InnoDB`, this value is an approximation, and may vary from the actual value by as much as 40 to 50%. In such cases, use `SELECT COUNT(*)` to obtain an accurate count.

  The `Rows` value is `NULL` for tables in the `INFORMATION_SCHEMA` database.

- `Avg_row_length`

  The average row length.

- `Data_length`

  For `MyISAM`, `Data_length` is the length of the data file, in bytes.

  For `InnoDB`, `Data_length` is the approximate amount of memory allocated for the clustered index, in bytes. Specifically, it is the clustered index size, in pages, multiplied by the `InnoDB` page size.

  Refer to the notes at the end of this section for information regarding other storage engines.

- `Max_data_length`

  For `MyISAM`, `Max_data_length` is maximum length of the data file. This is the total number of bytes of data that can be stored in the table, given the data pointer size used.

  Unused for `InnoDB`.

  Refer to the notes at the end of this section for information regarding other storage engines.

- `Index_length`

  For `MyISAM`, `Index_length` is the length of the index file, in bytes.

  For `InnoDB`, `Index_length` is the approximate amount of memory allocated for non-clustered indexes, in bytes. Specifically, it is the sum of non-clustered index sizes, in pages, multiplied by the `InnoDB` page size.

  Refer to the notes at the end of this section for information regarding other storage engines.

- `Data_free`

  The number of allocated but unused bytes.

  This information is also shown for `InnoDB` tables (previously, it was in the `Comment` value). `InnoDB` tables report the free space of the tablespace to which the table belongs. For a table located in the shared tablespace, this is the free space of the shared tablespace. If you are using multiple tablespaces and the table has its own tablespace, the free space is for only that table. Free space means the number of bytes in completely free extents minus a safety margin. Even if free space displays as 0, it may be possible to insert rows as long as new extents need not be allocated.

  For partitioned tables, this value is only an estimate and may not be absolutely correct. A more accurate method of obtaining this information in such cases is to query the `INFORMATION_SCHEMA.PARTITIONS` table, as shown in this example:

```
SELECT    SUM(DATA_FREE)
    FROM  INFORMATION_SCHEMA.PARTITIONS
```

```
    WHERE TABLE_SCHEMA = 'mydb'
    AND   TABLE_NAME   = 'mytable';
```

For more information, see Section 24.13, “The INFORMATION_SCHEMA PARTITIONS Table”.

- `Auto_increment`

  The next `AUTO_INCREMENT` value.

- `Create_time`

  When the table was created.

- `Update_time`

  When the data file was last updated. For some storage engines, this value is `NULL`. For example, `InnoDB` stores multiple tables in its system tablespace and the data file timestamp does not apply. Even with file-per-table mode with each `InnoDB` table in a separate `.ibd` file, change buffering can delay the write to the data file, so the file modification time is different from the time of the last insert, update, or delete. For `MyISAM`, the data file timestamp is used; however, on Windows the timestamp is not updated by updates so the value is inaccurate.

- `Check_time`

  When the table was last checked. Not all storage engines update this time, in which case the value is always `NULL`.

- `Collation`

  The table's character set and collation.

- `Checksum`

  The live checksum value (if any).

- `Create_options`

  Extra options used with `CREATE TABLE`. The original options supplied when `CREATE TABLE` is called are retained and the options reported here may differ from the active table settings and options.

- `Comment`

  The comment used when creating the table (or information as to why MySQL could not access the table information).

Notes:

- For `MEMORY` tables, the `Data_length`, `Max_data_length`, and `Index_length` values approximate the actual amount of allocated memory. The allocation algorithm reserves memory in large amounts to reduce the number of allocation operations.
- For views, all the fields displayed by `SHOW TABLE STATUS` are `NULL` except that `Name` indicates the view name and `Comment` says `view`.

13.7.5.37 SHOW TABLES Syntax

```
SHOW [EXTENDED] [FULL] TABLES
    [{FROM | IN} db_name]
    [LIKE 'pattern' | WHERE expr]
```

`SHOW TABLES` lists the non-`TEMPORARY` tables in a given database. You can also get this list using the `mysqlshow db_name` command. The `LIKE` clause, if present, indicates which table names to

match. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

Matching performed by the `LIKE` clause is dependent on the setting of the `lower_case_table_names` system variable.

As of MySQL 8.0.2, the optional `EXTENDED` modifier causes `SHOW TABLES` to list hidden tables created by failed `ALTER TABLE` statements. These temporary tables have names beginning with `#sql` and can be dropped using `DROP TABLE`.

This statement also lists any views in the database. The optional `FULL` modifier causes `SHOW TABLES` to display a second output column with values of `BASE TABLE` for a table and `VIEW` for a view.

If you have no privileges for a base table or view, it does not show up in the output from `SHOW TABLES` or `mysqlshow db_name`.

13.7.5.38 SHOW TRIGGERS Syntax

```
SHOW TRIGGERS
    [{FROM | IN} db_name]
    [LIKE 'pattern' | WHERE expr]
```

`SHOW TRIGGERS` lists the triggers currently defined for tables in a database (the default database unless a `FROM` clause is given). This statement returns results only for databases and tables for which you have the `TRIGGER` privilege. The `LIKE` clause, if present, indicates which table names to match (not trigger names) and causes the statement to display triggers for those tables. The `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, “Extensions to SHOW Statements”.

For the trigger `ins_sum` as defined in Section 23.3, “Using Triggers”, the output of this statement is as shown here:

```
mysql> SHOW TRIGGERS LIKE 'acc%'\G
*************************** 1. row ***************************
             Trigger: ins_sum
               Event: INSERT
               Table: account
           Statement: SET @sum = @sum + NEW.amount
              Timing: BEFORE
             Created: 2013-07-09 10:39:34.96
            sql_mode: NO_ENGINE_SUBSTITUTION
             Definer: me@localhost
character_set_client: utf8
collation_connection: utf8_general_ci
  Database Collation: utf8mb4_0900_ai_ci
```

`SHOW TRIGGERS` output has the following columns:

- `Trigger`: The trigger name.
- `Event`: The type of operation that causes trigger activation. The value is `'INSERT'`, `'UPDATE'`, or `'DELETE'`.
- `Table`: The table for which the trigger is defined.
- `Statement`: The trigger body; that is, the statement executed when the trigger activates.
- `Timing`: Whether the trigger activates before or after the triggering event. The value is `'BEFORE'` or `'AFTER'`.
- `Created`: The date and time when the trigger was created. This is a `TIMESTAMP(2)` value (with a fractional part in hundredths of seconds) for triggers.
- `sql_mode`: The SQL mode in effect when the trigger executes.

- `Definer`: The account of the user who created the trigger, in `'user_name'@'host_name'` format.
- `character_set_client`: The session value of the `character_set_client` system variable when the trigger was created.
- `collation_connection`: The session value of the `collation_connection` system variable when the trigger was created.
- `Database Collation`: The collation of the database with which the trigger is associated.

You can also obtain information about trigger objects from `INFORMATION_SCHEMA`, which contains a `TRIGGERS` table. See Section 24.28, "The INFORMATION_SCHEMA TRIGGERS Table".

13.7.5.39 SHOW VARIABLES Syntax

```
SHOW [GLOBAL | SESSION] VARIABLES
    [LIKE 'pattern' | WHERE expr]
```

`SHOW VARIABLES` shows the values of MySQL system variables (see Section 5.1.5, "Server System Variables"). This statement does not require any privilege. It requires only the ability to connect to the server.

System variable information is also available from these sources:

- Performance Schema tables. See Section 25.11.13, "Performance Schema System Variable Tables".
- The `mysqladmin variables` command. See Section 4.5.2, "`mysqladmin` — Client for Administering a MySQL Server".

For `SHOW VARIABLES`, a `LIKE` clause, if present, indicates which variable names to match. A `WHERE` clause can be given to select rows using more general conditions, as discussed in Section 24.33, "Extensions to SHOW Statements".

`SHOW VARIABLES` accepts an optional `GLOBAL` or `SESSION` variable scope modifier:

- With a `GLOBAL` modifier, the statement displays global system variable values. These are the values used to initialize the corresponding session variables for new connections to MySQL. If a variable has no global value, no value is displayed.
- With a `SESSION` modifier, the statement displays the system varaible values that are in effect for the current connection. If a variable has no session value, the global value is displayed. `LOCAL` is a synonym for `SESSION`.
- If no modifier is present, the default is `SESSION`.

The scope for each system variable is listed at Section 5.1.5, "Server System Variables".

`SHOW VARIABLES` is subject to a version-dependent display-width limit. For variables with very long values that are not completely displayed, use `SELECT` as a workaround. For example:

```
SELECT @@GLOBAL.innodb_data_file_path;
```

Most system variables can be set at server startup (read-only variables such as `version_comment` are exceptions). Many can be changed at runtime with the `SET` statement. See Section 5.1.6, "Using System Variables", and Section 13.7.4.1, "SET Syntax for Variable Assignment".

Partial output is shown here. The list of names and values may differ for your server. Section 5.1.5, "Server System Variables", describes the meaning of each variable, and Section 5.1.1, "Configuring the Server", provides information about tuning them.

```
mysql> SHOW VARIABLES;
+---------------------------------------+----------------------------+
| Variable_name                         | Value                      |
+---------------------------------------+----------------------------+
| auto_increment_increment              | 1                          |
| auto_increment_offset                 | 1                          |
| autocommit                            | ON                         |
| automatic_sp_privileges               | ON                         |
| back_log                              | 50                         |
| basedir                               | /home/jon/bin/mysql-5.5    |
| big_tables                            | OFF                        |
| binlog_cache_size                     | 32768                      |
| binlog_direct_non_transactional_updates | OFF                      |
| binlog_format                         | STATEMENT                  |
| binlog_stmt_cache_size                | 32768                      |
| bulk_insert_buffer_size               | 8388608                    |
...
| max_allowed_packet                    | 4194304                    |
| max_binlog_cache_size                 | 18446744073709547520       |
| max_binlog_size                       | 1073741824                 |
| max_binlog_stmt_cache_size            | 18446744073709547520       |
| max_connect_errors                    | 100                        |
| max_connections                       | 151                        |
| max_delayed_threads                   | 20                         |
| max_error_count                       | 1024                       |
| max_heap_table_size                   | 16777216                   |
| max_insert_delayed_threads            | 20                         |
| max_join_size                         | 18446744073709551615       |
...

| thread_handling                       | one-thread-per-connection  |
| thread_stack                          | 262144                     |
| time_format                           | %H:%i:%s                   |
| time_zone                             | SYSTEM                     |
| timestamp                             | 1316689732                 |
| tmp_table_size                        | 16777216                   |
| tmpdir                                | /tmp                       |
| transaction_alloc_block_size          | 8192                       |
| transaction_prealloc_size             | 4096                       |
| tx_isolation                          | REPEATABLE-READ            |
| unique_checks                         | ON                         |
| updatable_views_with_limit            | YES                        |
| version                               | 5.5.17-log                 |
| version_comment                       | Source distribution        |
| version_compile_machine               | x86_64                     |
| version_compile_os                    | Linux                      |
| wait_timeout                          | 28800                      |
| warning_count                         | 0                          |
+---------------------------------------+----------------------------+
```

With a `LIKE` clause, the statement displays only rows for those variables with names that match the pattern. To obtain the row for a specific variable, use a `LIKE` clause as shown:

```
SHOW VARIABLES LIKE 'max_join_size';
SHOW SESSION VARIABLES LIKE 'max_join_size';
```

To get a list of variables whose name match a pattern, use the `%` wildcard character in a `LIKE` clause:

```
SHOW VARIABLES LIKE '%size%';
SHOW GLOBAL VARIABLES LIKE '%size%';
```

Wildcard characters can be used in any position within the pattern to be matched. Strictly speaking, because `_` is a wildcard that matches any single character, you should escape it as `\_` to match it literally. In practice, this is rarely necessary.

13.7.5.40 SHOW WARNINGS Syntax

```
```

```
SHOW WARNINGS [LIMIT [offset,] row_count]
SHOW COUNT(*) WARNINGS
```

`SHOW WARNINGS` is a diagnostic statement that displays information about the conditions (errors, warnings, and notes) resulting from executing a statement in the current session. Warnings are generated for DML statements such as `INSERT`, `UPDATE`, and `LOAD DATA INFILE` as well as DDL statements such as `CREATE TABLE` and `ALTER TABLE`.

The `LIMIT` clause has the same syntax as for the `SELECT` statement. See Section 13.2.10, “SELECT Syntax”.

`SHOW WARNINGS` is also used following `EXPLAIN`, to display the extended information generated by `EXPLAIN`. See Section 8.8.3, “Extended EXPLAIN Output Format”.

`SHOW WARNINGS` displays information about the conditions resulting from execution of the most recent nondiagnostic statement in the current session. If the most recent statement resulted in an error during parsing, `SHOW WARNINGS` shows the resulting conditions, regardless of statement type (diagnostic or nondiagnostic).

The `SHOW COUNT(*) WARNINGS` diagnostic statement displays the total number of errors, warnings, and notes. You can also retrieve this number from the `warning_count` system variable:

```
SHOW COUNT(*) WARNINGS;
SELECT @@warning_count;
```

A difference in these statements is that the first is a diagnostic statement that does not clear the message list. The second, because it is a `SELECT` statement is considered nondiagnostic and does clear the message list.

A related diagnostic statement, `SHOW ERRORS`, shows only error conditions (it excludes warnings and notes), and `SHOW COUNT(*) ERRORS` statement displays the total number of errors. See Section 13.7.5.17, “SHOW ERRORS Syntax”. `GET DIAGNOSTICS` can be used to examine information for individual conditions. See Section 13.6.7.3, “GET DIAGNOSTICS Syntax”.

Here is a simple example that shows data-conversion warnings for `INSERT`:

```
mysql> CREATE TABLE t1 (a TINYINT NOT NULL, b CHAR(4));
Query OK, 0 rows affected (0.05 sec)

mysql> INSERT INTO t1 VALUES(10,'mysql'), (NULL,'test'), (300,'xyz');
Query OK, 3 rows affected, 3 warnings (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 3

mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1265
Message: Data truncated for column 'b' at row 1
*************************** 2. row ***************************
  Level: Warning
   Code: 1048
Message: Column 'a' cannot be null
*************************** 3. row ***************************
  Level: Warning
   Code: 1264
Message: Out of range value for column 'a' at row 3
3 rows in set (0.00 sec)
```

The `max_error_count` system variable controls the maximum number of error, warning, and note messages for which the server stores information, and thus the number of messages that `SHOW WARNINGS` displays. To change the number of messages the server can store, change the value of `max_error_count`.

`max_error_count` controls only how many messages are stored, not how many are counted. The value of `warning_count` is not limited by `max_error_count`, even if the number of messages generated exceeds `max_error_count`. The following example demonstrates this. The `ALTER TABLE` statement produces three warning messages (strict SQL mode is disabled for the example to prevent an error from occuring after a single conversion issue). Only one message is stored and displayed because `max_error_count` has been set to 1, but all three are counted (as shown by the value of `warning_count`):

```
mysql> SHOW VARIABLES LIKE 'max_error_count';
+-----------------+-------+
| Variable_name   | Value |
+-----------------+-------+
| max_error_count | 1024  |
+-----------------+-------+
1 row in set (0.00 sec)

mysql> SET max_error_count=1, sql_mode = '';
Query OK, 0 rows affected (0.00 sec)

mysql> ALTER TABLE t1 MODIFY b CHAR;
Query OK, 3 rows affected, 3 warnings (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 3

mysql> SHOW WARNINGS;
+---------+------+---------------------------------------+
| Level   | Code | Message                               |
+---------+------+---------------------------------------+
| Warning | 1263 | Data truncated for column 'b' at row 1 |
+---------+------+---------------------------------------+
1 row in set (0.00 sec)

mysql> SELECT @@warning_count;
+-----------------+
| @@warning_count |
+-----------------+
|               3 |
+-----------------+
1 row in set (0.01 sec)
```

To disable message storage, set `max_error_count` to 0. In this case, `warning_count` still indicates how many warnings occurred, but messages are not stored and cannot be displayed.

The `sql_notes` system variable controls whether note messages increment `warning_count` and whether the server stores them. By default, `sql_notes` is 1, but if set to 0, notes do not increment `warning_count` and the server does not store them:

```
mysql> SET sql_notes = 1;
mysql> DROP TABLE IF EXISTS test.no_such_table;
Query OK, 0 rows affected, 1 warning (0.00 sec)
mysql> SHOW WARNINGS;
+-------+------+-----------------------------------+
| Level | Code | Message                           |
+-------+------+-----------------------------------+
| Note  | 1051 | Unknown table 'test.no_such_table' |
+-------+------+-----------------------------------+
1 row in set (0.00 sec)

mysql> SET sql_notes = 0;
mysql> DROP TABLE IF EXISTS test.no_such_table;
Query OK, 0 rows affected (0.00 sec)
mysql> SHOW WARNINGS;
Empty set (0.00 sec)
```

The MySQL server sends to each client a count indicating the total number of errors, warnings, and notes resulting from the most recent statement executed by that client. From the C API, this value can be obtained by calling `mysql_warning_count()`. See Section 27.7.7.80, “mysql_warning_count()”.

In the `mysql` client, you can enable and disable automatic warnings display using the `warnings` and `nowarning` commands, respectively, or their shortcuts, `\W` and `\w` (see Section 4.5.1.2, “mysql Commands”). For example:

```
mysql> \W
Show warnings enabled.
mysql> SELECT 1/0;
+------+
| 1/0  |
+------+
| NULL |
+------+
1 row in set, 1 warning (0.03 sec)

Warning (Code 1365): Division by 0
mysql> \w
Show warnings disabled.
```

13.7.6 Other Administrative Statements

13.7.6.1 BINLOG Syntax

```
BINLOG 'str'
```

`BINLOG` is an internal-use statement. It is generated by the `mysqlbinlog` program as the printable representation of certain events in binary log files. (See Section 4.6.8, “`mysqlbinlog` — Utility for Processing Binary Log Files”.) The *`'str'`* value is a base 64-encoded string the that server decodes to determine the data change indicated by the corresponding event. This statement requires the `BINLOG_ADMIN` or `SUPER` privilege.

This statement can execute only format description events and row events.

13.7.6.2 CACHE INDEX Syntax

```
CACHE INDEX
  tbl_index_list [, tbl_index_list] ...
  [PARTITION (partition_list | ALL)]
  IN key_cache_name

tbl_index_list:
  tbl_name [[INDEX|KEY] (index_name[, index_name] ...)]

partition_list:
  partition_name[, partition_name][, ...]
```

The `CACHE INDEX` statement assigns table indexes to a specific key cache. It is used only for `MyISAM` tables. After the indexes have been assigned, they can be preloaded into the cache if desired with `LOAD INDEX INTO CACHE`.

The following statement assigns indexes from the tables `t1`, `t2`, and `t3` to the key cache named `hot_cache`:

```
mysql> CACHE INDEX t1, t2, t3 IN hot_cache;
+---------+--------------------+----------+----------+
| Table   | Op                 | Msg_type | Msg_text |
+---------+--------------------+----------+----------+
| test.t1 | assign_to_keycache | status   | OK       |
| test.t2 | assign_to_keycache | status   | OK       |
| test.t3 | assign_to_keycache | status   | OK       |
+---------+--------------------+----------+----------+
```

The syntax of `CACHE INDEX` enables you to specify that only particular indexes from a table should be assigned to the cache. The current implementation assigns all the table's indexes to the cache, so there is no reason to specify anything other than the table name.

The key cache referred to in a `CACHE INDEX` statement can be created by setting its size with a parameter setting statement or in the server parameter settings. For example:

```
mysql> SET GLOBAL keycache1.key_buffer_size=128*1024;
```

Key cache parameters can be accessed as members of a structured system variable. See Section 5.1.6.1, “Structured System Variables”.

A key cache must exist before you can assign indexes to it:

```
mysql> CACHE INDEX t1 IN non_existent_cache;
ERROR 1284 (HY000): Unknown key cache 'non_existent_cache'
```

By default, table indexes are assigned to the main (default) key cache created at the server startup. When a key cache is destroyed, all indexes assigned to it become assigned to the default key cache again.

Index assignment affects the server globally: If one client assigns an index to a given cache, this cache is used for all queries involving the index, no matter which client issues the queries.

In MySQL 8.0, this statement is also supported for partitioned `MyISAM` tables. You can assign one or more indexes for one, several, or all partitions to a given key cache. For example, you can do the following:

```
CREATE TABLE pt (c1 INT, c2 VARCHAR(50), INDEX i(c1))
    ENGINE=MyISAM
    PARTITION BY HASH(c1)
    PARTITIONS 4;

SET GLOBAL kc_fast.key_buffer_size = 128 * 1024;
SET GLOBAL kc_slow.key_buffer_size = 128 * 1024;

CACHE INDEX pt PARTITION (p0) IN kc_fast;
CACHE INDEX pt PARTITION (p1, p3) IN kc_slow;
```

The previous set of statements performs the following actions:

- Creates a partitioned table with 4 partitions; these partitions are automatically named `p0`, ..., `p3`; this table has an index named `i` on column `c1`.
- Creates 2 key caches named `kc_fast` and `kc_slow`
- Assigns the index for partition `p0` to the `kc_fast` key cache and the index for partitions `p1` and `p3` to the `kc_slow` key cache; the index for the remaining partition (`p2`) uses the server's default key cache.

If you wish instead to assign the indexes for all partitions in table `pt` to a single key cache named `kc_all`, you can use either one of the following 2 statements:

```
CACHE INDEX pt PARTITION (ALL) IN kc_all;

CACHE INDEX pt IN kc_all;
```

The two statements just shown are equivalent, and issuing either one of them has exactly the same effect. In other words, if you wish to assign indexes for all partitions of a partitioned table to the same key cache, then the `PARTITION (ALL)` clause is optional.

When assigning indexes for multiple partitions to a key cache, the partitions do not have to be contiguous, and you are not required to list their names in any particular order. Indexes for any partitions that are not explicitly assigned to a key cache automatically use the server's default key cache.

In MySQL 8.0, index preloading is also supported for partitioned MyISAM tables. For more information, see Section 13.7.6.5, “LOAD INDEX INTO CACHE Syntax”.

13.7.6.3 FLUSH Syntax

```
FLUSH [NO_WRITE_TO_BINLOG | LOCAL]
    flush_option [, flush_option] ...
```

The FLUSH statement has several variant forms that clear or reload various internal caches, flush tables, or acquire locks. To execute FLUSH, you must have the RELOAD privilege. Specific flush options might require additional privileges, as described later.

By default, the server writes FLUSH statements to the binary log so that they replicate to replication slaves. To suppress logging, specify the optional NO_WRITE_TO_BINLOG keyword or its alias LOCAL.

Note

FLUSH LOGS, FLUSH TABLES WITH READ LOCK (with or without a table list), and FLUSH TABLES *tbl_name* ... FOR EXPORT are not written to the binary log in any case because they would cause problems if replicated to a slave.

Sending a SIGHUP signal to the server causes several flush operations to occur that are similar to various forms of the FLUSH statement. See Section 5.1.11, “Server Response to Signals”.

The FLUSH statement causes an implicit commit. See Section 13.3.3, “Statements That Cause an Implicit Commit”.

The RESET statement is similar to FLUSH. See Section 13.7.6.6, “RESET Syntax”, for information about using the RESET statement with replication.

flush_option can be any of the following items.

- DES_KEY_FILE

  Reloads the DES keys from the file that was specified with the --des-key-file option at server startup time.

- HOSTS

  Empties the host cache. You should flush the host cache if some of your hosts change IP address or if the error message Host '*host_name*' is blocked occurs. (See Section B.5.2.5, “Host 'host_name' is blocked”.) When more than max_connect_errors errors occur successively for a given host while connecting to the MySQL server, MySQL assumes that something is wrong and blocks the host from further connection requests. Flushing the host cache enables further connection attempts from the host. The default value of max_connect_errors is 100. To avoid this error message, start the server with max_connect_errors set to a large value.

- [*log_type*] LOGS | RELAY LOGS [*channel_option*]

  With no *log_type* option, FLUSH LOGS closes and reopens all log files. If binary logging is enabled, the sequence number of the binary log file is incremented by one relative to the previous file.

  FLUSH LOGS has no effect on tables used for the general query log or for the slow query log (see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”).

  With a *log_type* option, only the specified log type is flushed. These *log_type* options are permitted:

  - BINARY closes and reopens the binary log files. If binary logging is enabled, the sequence number of the binary log file is incremented by one relative to the previous file.

- `ENGINE` closes and reopens any flushable logs for installed storage engines. This causes `InnoDB` to flush its logs to disk.
- `ERROR` closes and reopens the error log file.
- `GENERAL` closes and reopens the general query log file.
- `RELAY` closes and reopens the relay log files. If binary logging is enabled, the sequence number of the binary log file is incremented by one relative to the previous file.

  The `FOR CHANNEL channel` clause enables you to choose which replication channel to apply a `FLUSH RELAY LOGS` statement to. If no clause is set and no extra replication channels exist, the statement applies to the default channel and behaves the same as versions of MySQL prior to 5.7.6. If multiple replication channels exist and no clause is set, *all* replication channels are flushed. Execute a `FLUSH RELAY LOGS FOR CHANNEL channel` statement to flush a specific replication channel's relay log. See Section 18.2.3, “Replication Channels” for more information.

  Note

  The `FOR CHANNEL channel` clause can only be used with the `RELAY LOGS log_type`.

- `SLOW` closes and reopens the slow query log file.

- `OPTIMIZER_COSTS`

  Rereads the cost model tables so that the optimizer starts using the current cost estimates stored in them. The server writes a warning to the error log for any unrecognized entries. (For information about these tables, see Section 8.9.5, “The Optimizer Cost Model”.) This operation affects only sessions that begin subsequent to the flush. Existing sessions continue to use the cost estimates that were current when they began.

- `PRIVILEGES`

  Reloads the privileges from the grant tables in the `mysql` database.

  The `FLUSH PRIVILEGES` statement reads the `global_grants` table for dynamic privilege assignments and registers any unregistered privileges found there.

  The server caches information in memory as a result of `GRANT`, `CREATE USER`, `CREATE SERVER`, and `INSTALL PLUGIN` statements. This memory is not released by the corresponding `REVOKE`, `DROP USER`, `DROP SERVER`, and `UNINSTALL PLUGIN` statements, so for a server that executes many instances of the statements that cause caching, there will be an increase in memory use. This cached memory can be freed with `FLUSH PRIVILEGES`.

- `QUERY CACHE`

  Defragment the query cache to better utilize its memory. `FLUSH QUERY CACHE` does not remove any queries from the cache, unlike `FLUSH TABLES` or `RESET QUERY CACHE`.

- `STATUS`

  This option adds the session status from all active sessions to the global status variables, resets the status of all active sessions, and resets account, host, and user status values aggregated from disconnected sessions. This is something that may be of use when debugging a query. See Section 1.7, “How to Report Bugs or Problems”.

- `TABLES`

  `FLUSH TABLES` flushes tables, and, depending on the variant used, acquires locks. The permitted syntax is discussed later in this section.

- `USER_RESOURCES`

  Resets all per-hour user resources to zero. This enables clients that have reached their hourly connection, query, or update limits to resume activity immediately. `FLUSH USER_RESOURCES` does not apply to the limit on maximum simultaneous connections. See Section 6.3.6, “Setting Account Resource Limits”.

The `mysqladmin` utility provides a command-line interface to some flush operations, using commands such as `flush-hosts`, `flush-logs`, `flush-privileges`, `flush-status`, and `flush-tables`. See Section 4.5.2, “`mysqladmin` — Client for Administering a MySQL Server”.

Note

It is not possible to issue `FLUSH` statements within stored functions or triggers. However, you may use `FLUSH` in stored procedures, so long as these are not called from stored functions or triggers. See Section C.1, “Restrictions on Stored Programs”.

FLUSH TABLES Syntax

`FLUSH TABLES` has several forms, described following. If any variant of the `TABLES` option is used in a `FLUSH` statement, it must be the only option used. `FLUSH TABLE` is a synonym for `FLUSH TABLES`.

- `FLUSH TABLES`

  Closes all open tables, forces all tables in use to be closed, and flushes the query cache. `FLUSH TABLES` also removes all query results from the query cache, like the `RESET QUERY CACHE` statement.

  In MySQL 8.0, `FLUSH TABLES` is not permitted when there is an active `LOCK TABLES ... READ`. To flush and lock tables, use `FLUSH TABLES *tbl_name* ... WITH READ LOCK` instead.

- `FLUSH TABLES *tbl_name* [, *tbl_name*] ...`

  With a list of one or more comma-separated table names, this statement is like `FLUSH TABLES` with no names except that the server flushes only the named tables. No error occurs if a named table does not exist.

- `FLUSH TABLES WITH READ LOCK`

  Closes all open tables and locks all tables for all databases with a global read lock. This is a very convenient way to get backups if you have a file system such as Veritas or ZFS that can take snapshots in time. Use `UNLOCK TABLES` to release the lock.

  `FLUSH TABLES WITH READ LOCK` acquires a global read lock and not table locks, so it is not subject to the same behavior as `LOCK TABLES` and `UNLOCK TABLES` with respect to table locking and implicit commits:

  - `UNLOCK TABLES` implicitly commits any active transaction only if any tables currently have been locked with `LOCK TABLES`. The commit does not occur for `UNLOCK TABLES` following `FLUSH TABLES WITH READ LOCK` because the latter statement does not acquire table locks.
  - Beginning a transaction causes table locks acquired with `LOCK TABLES` to be released, as though you had executed `UNLOCK TABLES`. Beginning a transaction does not release a global read lock acquired with `FLUSH TABLES WITH READ LOCK`.

  `FLUSH TABLES WITH READ LOCK` does not prevent the server from inserting rows into the log tables (see Section 5.4.1, “Selecting General Query and Slow Query Log Output Destinations”).

  `FLUSH TABLES WITH READ LOCK` is not compatible with XA transactions.

- `FLUSH TABLES *tbl_name* [, *tbl_name*] ... WITH READ LOCK`

This statement flushes and acquires read locks for the named tables. The statement first acquires exclusive metadata locks for the tables, so it waits for transactions that have those tables open to complete. Then the statement flushes the tables from the table cache, reopens the tables, acquires table locks (like `LOCK TABLES ... READ`), and downgrades the metadata locks from exclusive to shared. After the statement acquires locks and downgrades the metadata locks, other sessions can read but not modify the tables.

Because this statement acquires table locks, you must have the `LOCK TABLES` privilege for each table, in addition to the `RELOAD` privilege that is required to use any `FLUSH` statement.

This statement applies only to existing base tables. If a name refers to a base table, that table is used. If it refers to a `TEMPORARY` table, it is ignored. If a name applies to a view, an `ER_WRONG_OBJECT` error occurs. Otherwise, an `ER_NO_SUCH_TABLE` error occurs.

Use `UNLOCK TABLES` to release the locks, `LOCK TABLES` to release the locks and acquire other locks, or `START TRANSACTION` to release the locks and begin a new transaction.

This variant of `FLUSH` enables tables to be flushed and locked in a single operation. It provides a workaround for the restriction in MySQL 8.0 that `FLUSH TABLES` is not permitted when there is an active `LOCK TABLES ... READ`.

This statement does not perform an implicit `UNLOCK TABLES`, so an error results if you use the statement while there is any active `LOCK TABLES` or use it a second time without first releasing the locks acquired.

If a flushed table was opened with `HANDLER`, the handler is implicitly flushed and loses its position.

- `FLUSH TABLES` *`tbl_name`* `[,` *`tbl_name`*`] ... FOR EXPORT`

  This `FLUSH TABLES` variant applies to `InnoDB` tables. It ensures that changes to the named tables have been flushed to disk so that binary table copies can be made while the server is running.

  The statement works like this:

  1. It acquires shared metadata locks for the named tables. The statement blocks as long as other sessions have active transactions that have modified those tables or hold table locks for them. When the locks have been acquired, the statement blocks transactions that attempt to update the tables while permitting read-only operations to continue.
  2. It checks whether all storage engines for the tables support `FOR EXPORT`. If any do not, an `ER_ILLEGAL_HA` error occurs and the statement fails.
  3. The statement notifies the storage engine for each table to make the table ready for export. The storage engine must ensure that any pending changes are written to disk.
  4. The statement puts the session in lock-tables mode so that the metadata locks acquired earlier are not released when the `FOR EXPORT` statement completes.

  The `FLUSH TABLES ... FOR EXPORT` statement requires that you have the `SELECT` privilege for each table. Because this statement acquires table locks, you must also have the `LOCK TABLES` privilege for each table, in addition to the `RELOAD` privilege that is required to use any `FLUSH` statement.

  This statement applies only to existing base tables. If a name refers to a base table, that table is used. If it refers to a `TEMPORARY` table, it is ignored. If a name applies to a view, an `ER_WRONG_OBJECT` error occurs. Otherwise, an `ER_NO_SUCH_TABLE` error occurs.

  `InnoDB` supports `FOR EXPORT` for tables that have their own .ibd file (that is, tables that were created with the `innodb_file_per_table` setting enabled). `InnoDB` ensures when notified by the `FOR EXPORT` statement that any changes have been flushed to disk. This permits a binary copy

of table contents to be made while the FOR EXPORT statement is in effect because the .ibd file is transaction consistent and can be copied while the server is running. FOR EXPORT does not apply to InnoDB system tablespace files, or to InnoDB tables that have any FULLTEXT indexes.

FLUSH TABLES ...FOR EXPORT is supported for partitioned InnoDB tables.

When notified by FOR EXPORT, InnoDB writes to disk certain kinds of data that is normally held in memory or in separate disk buffers outside the tablespace files. For each table, InnoDB also produces a file named *table_name*.cfg in the same database directory as the table. The .cfg file contains metadata needed to reimport the tablespace files later, into the same or different server.

When the FOR EXPORT statement completes, InnoDB will have flushed all dirty pages to the table data files. Any change buffer entries are merged prior to flushing. At this point, the tables are locked and quiescent: The tables are in a transactionally consistent state on disk and you can copy the .ibd tablespace files along with the corresponding .cfg files to get a consistent snapshot of those tables.

For the procedure to reimport the copied table data into a MySQL instance, see Section 15.7.6, "Copying File-Per-Table Tablespaces to Another Instance".

After you are done with the tables, use UNLOCK TABLES to release the locks, LOCK TABLES to release the locks and acquire other locks, or START TRANSACTION to release the locks and begin a new transaction.

While any of these statements is in effect within the session, attempts to use FLUSH TABLES ... FOR EXPORT produce an error:

```
FLUSH TABLES ... WITH READ LOCK
FLUSH TABLES ... FOR EXPORT
LOCK TABLES ... READ
LOCK TABLES ... WRITE
```

While FLUSH TABLES ... FOR EXPORT is in effect within the session, attempts to use any of these statements produce an error:

```
FLUSH TABLES WITH READ LOCK
FLUSH TABLES ... WITH READ LOCK
FLUSH TABLES ... FOR EXPORT
```

13.7.6.4 KILL Syntax

```
KILL [CONNECTION | QUERY] processlist_id
```

Each connection to mysqld runs in a separate thread. You can kill a thread with the KILL *processlist_id* statement.

Thread processlist identifiers can be determined from the ID column of the INFORMATION_SCHEMA.PROCESSLIST table, the Id column of SHOW PROCESSLIST output, and the PROCESSLIST_ID column of the Performance Schema threads table. The value for the current thread is returned by the CONNECTION_ID() function.

KILL permits an optional CONNECTION or QUERY modifier:

- KILL CONNECTION is the same as KILL with no modifier: It terminates the connection associated with the given *processlist_id*, after terminating any statement the connection is executing.
- KILL QUERY terminates the statement the connection is currently executing, but leaves the connection itself intact.

If you have the PROCESS privilege, you can see all threads. If you have the SUPER privilege, you can kill all threads and statements. Otherwise, you can see and kill only your own threads and statements.

You can also use the `mysqladmin processlist` and `mysqladmin kill` commands to examine and kill threads.

When you use `KILL`, a thread-specific kill flag is set for the thread. In most cases, it might take some time for the thread to die because the kill flag is checked only at specific intervals:

- During `SELECT` operations, for `ORDER BY` and `GROUP BY` loops, the flag is checked after reading a block of rows. If the kill flag is set, the statement is aborted.
- `ALTER TABLE` operations that make a table copy check the kill flag periodically for each few copied rows read from the original table. If the kill flag was set, the statement is aborted and the temporary table is deleted.

  The `KILL` statement returns without waiting for confirmation, but the kill flag check aborts the operation within a reasonably small amount of time. Aborting the operation to perform any necessary cleanup also takes some time.
- During `UPDATE` or `DELETE` operations, the kill flag is checked after each block read and after each updated or deleted row. If the kill flag is set, the statement is aborted. If you are not using transactions, the changes are not rolled back.
- `GET_LOCK()` aborts and returns `NULL`.
- If the thread is in the table lock handler (state: `Locked`), the table lock is quickly aborted.
- If the thread is waiting for free disk space in a write call, the write is aborted with a “disk full” error message.

Warning

Killing a `REPAIR TABLE` or `OPTIMIZE TABLE` operation on a `MyISAM` table results in a table that is corrupted and unusable. Any reads or writes to such a table fail until you optimize or repair it again (without interruption).

13.7.6.5 LOAD INDEX INTO CACHE Syntax

```
LOAD INDEX INTO CACHE
  tbl_index_list [, tbl_index_list] ...

tbl_index_list:
  tbl_name
    [PARTITION (partition_list | ALL)]
    [[INDEX|KEY] (index_name[, index_name] ...)]
    [IGNORE LEAVES]

partition_list:
    partition_name[, partition_name][, ...]
```

The `LOAD INDEX INTO CACHE` statement preloads a table index into the key cache to which it has been assigned by an explicit `CACHE INDEX` statement, or into the default key cache otherwise.

`LOAD INDEX INTO CACHE` is used only for `MyISAM` tables. In MySQL 8.0, it is also supported for partitioned `MyISAM` tables; in addition, indexes on partitioned tables can be preloaded for one, several, or all partitions.

The `IGNORE LEAVES` modifier causes only blocks for the nonleaf nodes of the index to be preloaded.

`IGNORE LEAVES` is also supported for partitioned `MyISAM` tables.

The following statement preloads nodes (index blocks) of indexes for the tables `t1` and `t2`:

```
mysql> LOAD INDEX INTO CACHE t1, t2 IGNORE LEAVES;
+---------+--------------+----------+----------+
```

```
| Table   | Op           | Msg_type | Msg_text |
+---------+--------------+----------+----------+
| test.t1 | preload_keys | status   | OK       |
| test.t2 | preload_keys | status   | OK       |
+---------+--------------+----------+----------+
```

This statement preloads all index blocks from `t1`. It preloads only blocks for the nonleaf nodes from `t2`.

The syntax of `LOAD INDEX INTO CACHE` enables you to specify that only particular indexes from a table should be preloaded. The current implementation preloads all the table's indexes into the cache, so there is no reason to specify anything other than the table name.

It is possible to preload indexes on specific partitions of partitioned `MyISAM` tables. For example, of the following 2 statements, the first preloads indexes for partition `p0` of a partitioned table `pt`, while the second preloads the indexes for partitions `p1` and `p3` of the same table:

```
LOAD INDEX INTO CACHE pt PARTITION (p0);
LOAD INDEX INTO CACHE pt PARTITION (p1, p3);
```

To preload the indexes for all partitions in table `pt`, you can use either one of the following 2 statements:

```
LOAD INDEX INTO CACHE pt PARTITION (ALL);

LOAD INDEX INTO CACHE pt;
```

The two statements just shown are equivalent, and issuing either one of them has exactly the same effect. In other words, if you wish to preload indexes for all partitions of a partitioned table, then the `PARTITION (ALL)` clause is optional.

When preloading indexes for multiple partitions, the partitions do not have to be contiguous, and you are not required to list their names in any particular order.

`LOAD INDEX INTO CACHE ... IGNORE LEAVES` fails unless all indexes in a table have the same block size. You can determine index block sizes for a table by using `myisamchk -dv` and checking the `Blocksize` column.

13.7.6.6 RESET Syntax

```
RESET reset_option [, reset_option] ...
```

The `RESET` statement is used to clear the state of various server operations. You must have the `RELOAD` privilege to execute `RESET`.

For information about the `RESET PERSIST` statement that removes persisted global system variables, see Section 13.7.6.7, “RESET PERSIST Syntax”.

`RESET` acts as a stronger version of the `FLUSH` statement. See Section 13.7.6.3, “FLUSH Syntax”.

The `RESET` statement causes an implicit commit. See Section 13.3.3, “Statements That Cause an Implicit Commit”.

`reset_option` can be any of the following:

- `MASTER`

  Deletes all binary logs listed in the index file, resets the binary log index file to be empty, and creates a new binary log file.

- `QUERY CACHE`

  Removes all query results from the query cache.

- `SLAVE`

  Makes the slave forget its replication position in the master binary logs. Also resets the relay log by deleting any existing relay log files and beginning a new one.

13.7.6.7 RESET PERSIST Syntax

```
RESET PERSIST [[IF EXISTS] system_var_name]
```

`RESET PERSIST` is used to remove persisted global system variable settings from the `mysqld-auto.cnf` file. This statement requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege. Removing a persisted system variable causes the variable no longer to be initialized from `mysqld-auto.cnf` at server startup.

For information about persisted system variables, see Section 5.1.6, “Using System Variables”. For information about `RESET` statement variants that clear the state of other server operations, see Section 13.7.6.6, “RESET Syntax”.

Depending on whether the variable name and `IF EXISTS` clauses are present, the statement has these forms:

- `RESET PERSIST`

  This statement removes all persisted variables from `mysqld-auto.cnf`.

- `RESET PERSIST system_var_name`

  This statement removes the named persisted variable from `mysqld-auto.cnf`. This includes plugin system variables, even if the plugin is not currently installed. If the variable is not present in the file, an error occurs.

- `RESET PERSIST IF EXISTS system_var_name`

  This statement removes the named persisted variable from `mysqld-auto.cnf`. This includes plugin system variables, even if the plugin is not currently installed. If the variable is not present in the file, a warning is generated.

`RESET PERSIST` is not affected by the value of the `persisted_globals_load` system variable.

`RESET PERSIST` affects the contents of the Performance Schema `persisted_variables` table because the table contents correspond to the contents of the `mysqld-auto.cnf` file.

`RESET PERSIST` has no effect on the contents of the Performance Schema `variables_info` table.

13.7.6.8 SHUTDOWN Syntax

```
SHUTDOWN
```

This statement stops the MySQL server. It requires the `SHUTDOWN` privilege.

`SHUTDOWN` provides an SQL-level interface to the same functionality available using the `mysqladmin shutdown` command.

13.8 Utility Statements

13.8.1 DESCRIBE Syntax

The `DESCRIBE` and `EXPLAIN` statements are synonyms, used either to obtain information about table structure or query execution plans. For more information, see Section 13.7.5.5, “SHOW COLUMNS Syntax”, and Section 13.8.2, “EXPLAIN Syntax”.

13.8.2 EXPLAIN Syntax

```
{EXPLAIN | DESCRIBE | DESC}
    tbl_name [col_name | wild]

{EXPLAIN | DESCRIBE | DESC}
    [explain_type]
    {explainable_stmt | FOR CONNECTION connection_id}

explain_type: {
    EXTENDED
  | PARTITIONS
  | FORMAT = format_name
}

format_name: {
    TRADITIONAL
  | JSON
}

explainable_stmt: {
    SELECT statement
  | DELETE statement
  | INSERT statement
  | REPLACE statement
  | UPDATE statement
}
```

The DESCRIBE and EXPLAIN statements are synonyms. In practice, the DESCRIBE keyword is more often used to obtain information about table structure, whereas EXPLAIN is used to obtain a query execution plan (that is, an explanation of how MySQL would execute a query).

The following discussion uses the DESCRIBE and EXPLAIN keywords in accordance with those uses, but the MySQL parser treats them as completely synonymous.

- Obtaining Table Structure Information
- Obtaining Execution Plan Information

Obtaining Table Structure Information

DESCRIBE provides information about the columns in a table:

```
mysql> DESCRIBE City;
+------------+----------+------+-----+---------+----------------+
| Field      | Type     | Null | Key | Default | Extra          |
+------------+----------+------+-----+---------+----------------+
| Id         | int(11)  | NO   | PRI | NULL    | auto_increment |
| Name       | char(35) | NO   |     |         |                |
| Country    | char(3)  | NO   | UNI |         |                |
| District   | char(20) | YES  | MUL |         |                |
| Population | int(11)  | NO   |     | 0       |                |
+------------+----------+------+-----+---------+----------------+
```

DESCRIBE is a shortcut for SHOW COLUMNS. These statements also display information for views. The description for SHOW COLUMNS provides more information about the output columns. See Section 13.7.5.5, “SHOW COLUMNS Syntax”.

By default, DESCRIBE displays information about all columns in the table. *col_name*, if given, is the name of a column in the table. In this case, the statement displays information only for the named column. *wild*, if given, is a pattern string. It can contain the SQL % and _ wildcard characters. In this case, the statement displays output only for the columns with names matching the string. There is no need to enclose the string within quotation marks unless it contains spaces or other special characters.

The DESCRIBE statement is provided for compatibility with Oracle.

The `SHOW CREATE TABLE`, `SHOW TABLE STATUS`, and `SHOW INDEX` statements also provide information about tables. See Section 13.7.5, “SHOW Syntax”.

Obtaining Execution Plan Information

The `EXPLAIN` statement provides information about how MySQL executes statements:

- `EXPLAIN` works with `SELECT`, `DELETE`, `INSERT`, `REPLACE`, and `UPDATE` statements.
- When `EXPLAIN` is used with an explainable statement, MySQL displays information from the optimizer about the statement execution plan. That is, MySQL explains how it would process the statement, including information about how tables are joined and in which order. For information about using `EXPLAIN` to obtain execution plan information, see Section 8.8.2, “EXPLAIN Output Format”.
- When `EXPLAIN` is used with `FOR CONNECTION` *`connection_id`* rather than an explainable statement, it displays the execution plan for the statement executing in the named connection. See Section 8.8.4, “Obtaining Execution Plan Information for a Named Connection”.
- For `SELECT` statements, `EXPLAIN` produces additional execution plan information. See Section 8.8.3, “Extended EXPLAIN Output Format”.

Note

In older MySQL releases, extended information was produced using `EXPLAIN EXTENDED`. That syntax is still recognized for backward compatibility but extended output is now enabled by default, so the `EXTENDED` keyword is superfluous and deprecated. Its use results in a warning, and it will be removed from `EXPLAIN` syntax in a future MySQL release.

- `EXPLAIN` is useful for examining queries involving partitioned tables. See Section 22.3.5, “Obtaining Information About Partitions”.

Note

In older MySQL releases, partition information was produced using `EXPLAIN PARTITIONS`. That syntax is still recognized for backward compatibility but partition output is now enabled by default, so the `PARTITIONS` keyword is superfluous and deprecated. Its use results in a warning, and it will be removed from `EXPLAIN` syntax in a future MySQL release.

- The `FORMAT` option can be used to select the output format. `TRADITIONAL` presents the output in tabular format. This is the default if no `FORMAT` option is present. `JSON` format displays the information in JSON format.

`EXPLAIN` requires the `SELECT` privilege for any tables or views accessed, including any underlying tables of views. For views, `EXPLAIN` also requires the `SHOW VIEW` privilege.

With the help of `EXPLAIN`, you can see where you should add indexes to tables so that the statement executes faster by using indexes to find rows. You can also use `EXPLAIN` to check whether the optimizer joins the tables in an optimal order. To give a hint to the optimizer to use a join order corresponding to the order in which the tables are named in a `SELECT` statement, begin the statement with `SELECT STRAIGHT_JOIN` rather than just `SELECT`. (See Section 13.2.10, “SELECT Syntax”.)

The optimizer trace may sometimes provide information complementary to that of `EXPLAIN`. However, the optimizer trace format and content are subject to change between versions. For details, see MySQL Internals: Tracing the Optimizer.

If you have a problem with indexes not being used when you believe that they should be, run `ANALYZE TABLE` to update table statistics, such as cardinality of keys, that can affect the choices the optimizer makes. See Section 13.7.2.1, “ANALYZE TABLE Syntax”.

Note

MySQL Workbench has a Visual Explain capability that provides a visual representation of `EXPLAIN` output. See Tutorial: Using Visual Explain to Improve Query Performance.

13.8.3 HELP Syntax

```
HELP 'search_string'
```

The `HELP` statement returns online information from the MySQL Reference manual. Its proper operation requires that the help tables in the `mysql` database be initialized with help topic information (see Section 5.1.10, “Server-Side Help”).

The `HELP` statement searches the help tables for the given search string and displays the result of the search. The search string is not case sensitive.

The search string can contain the wildcard characters `%` and `_`. These have the same meaning as for pattern-matching operations performed with the `LIKE` operator. For example, `HELP 'rep%'` returns a list of topics that begin with `rep`.

The HELP statement understands several types of search strings:

- At the most general level, use `contents` to retrieve a list of the top-level help categories:

  ```
  HELP 'contents'
  ```

- For a list of topics in a given help category, such as `Data Types`, use the category name:

  ```
  HELP 'data types'
  ```

- For help on a specific help topic, such as the `ASCII()` function or the `CREATE TABLE` statement, use the associated keyword or keywords:

  ```
  HELP 'ascii'
  HELP 'create table'
  ```

In other words, the search string matches a category, many topics, or a single topic. You cannot necessarily tell in advance whether a given search string will return a list of items or the help information for a single help topic. However, you can tell what kind of response `HELP` returned by examining the number of rows and columns in the result set.

The following descriptions indicate the forms that the result set can take. Output for the example statements is shown using the familiar “tabular” or “vertical” format that you see when using the `mysql` client, but note that `mysql` itself reformats `HELP` result sets in a different way.

- Empty result set

  No match could be found for the search string.

- Result set containing a single row with three columns

  This means that the search string yielded a hit for the help topic. The result has three columns:

  - `name`: The topic name.
  - `description`: Descriptive help text for the topic.
  - `example`: Usage example or examples. This column might be blank.

  Example: `HELP 'replace'`

Yields:

```
name: REPLACE
description: Syntax:
REPLACE(str,from_str,to_str)

Returns the string str with all occurrences of the string from_str
replaced by the string to_str. REPLACE() performs a case-sensitive
match when searching for from_str.
example: mysql> SELECT REPLACE('www.mysql.com', 'w', 'Ww');
        -> 'WwWwWw.mysql.com'
```

- Result set containing multiple rows with two columns

  This means that the search string matched many help topics. The result set indicates the help topic names:

  - `name`: The help topic name.
  - `is_it_category`: `Y` if the name represents a help category, `N` if it does not. If it does not, the `name` value when specified as the argument to the `HELP` statement should yield a single-row result set containing a description for the named item.

  Example: `HELP 'status'`

  Yields:

```
+----------------------+----------------+
| name                 | is_it_category |
+----------------------+----------------+
| SHOW                 | N              |
| SHOW ENGINE          | N              |
| SHOW MASTER STATUS   | N              |
| SHOW PROCEDURE STATUS | N              |
| SHOW SLAVE STATUS    | N              |
| SHOW STATUS          | N              |
| SHOW TABLE STATUS    | N              |
+----------------------+----------------+
```

- Result set containing multiple rows with three columns

  This means the search string matches a category. The result set contains category entries:

  - `source_category_name`: The help category name.
  - `name`: The category or topic name
  - `is_it_category`: `Y` if the name represents a help category, `N` if it does not. If it does not, the `name` value when specified as the argument to the `HELP` statement should yield a single-row result set containing a description for the named item.

  Example: `HELP 'functions'`

  Yields:

```
+---------------------+-------------------------+----------------+
| source_category_name | name                    | is_it_category |
+---------------------+-------------------------+----------------+
| Functions           | CREATE FUNCTION         | N              |
| Functions           | DROP FUNCTION           | N              |
| Functions           | Bit Functions           | Y              |
| Functions           | Comparison operators    | Y              |
| Functions           | Control flow functions  | Y              |
```

```
| Functions            | Date and Time Functions | Y              |
| Functions            | Encryption Functions    | Y              |
| Functions            | Information Functions   | Y              |
| Functions            | Logical operators       | Y              |
| Functions            | Miscellaneous Functions | Y              |
| Functions            | Numeric Functions       | Y              |
| Functions            | String Functions        | Y              |
+----------------------+-------------------------+----------------+
```

13.8.4 USE Syntax

```
USE db_name
```

The `USE db_name` statement tells MySQL to use the *`db_name`* database as the default (current) database for subsequent statements. The database remains the default until the end of the session or another `USE` statement is issued:

```
USE db1;
SELECT COUNT(*) FROM mytable;   # selects from db1.mytable
USE db2;
SELECT COUNT(*) FROM mytable;   # selects from db2.mytable
```

Making a particular database the default by means of the `USE` statement does not preclude you from accessing tables in other databases. The following example accesses the `author` table from the `db1` database and the `editor` table from the `db2` database:

```
USE db1;
SELECT author_name,editor_name FROM author,db2.editor
  WHERE author.editor_id = db2.editor.editor_id;
```

Chapter 14 MySQL Data Dictionary

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MySQL Server now incorporates a transactional data dictionary that stores information about database objects. In previous MySQL releases, dictionary data was stored in metadata files and nontransactional tables.

This chapter describes the main features, benefits, usage differences, and limitations of the data dictionary. For other implications of the data dictionary feature, refer to the “Data Dictionary Notes” section in the MySQL 8.0 Release Notes.

In in MySQL 8.0.0 and MySQL 8.0.1, the `InnoDB` storage engine uses its own storage engine-specific data dictionary. As of MySQL 8.0.2, `InnoDB` uses the MySQL data dictionary.

Benefits of the MySQL data dictionary include:

- Simplicity of a centralized data dictionary schema that uniformly stores dictionary data. See Section 14.1, “Data Dictionary Schema”.
- Removal of file-based metadata storage. See Section 14.2, “Removal of File-based Metadata Storage”.
- Transactional, crash-safe storage of dictionary data. See Section 14.3, “Transactional Storage of Dictionary Data”.
- Uniform and centralized caching for dictionary objects. See Section 14.4, “Dictionary Object Cache”.
- A simpler and improved implementation for some `INFORMATION_SCHEMA` tables. See Section 14.5, “INFORMATION_SCHEMA and Data Dictionary Integration”.

Important

A data dictionary-enabled server entails some general operational differences compared to a server that does not have a data dictionary; see Section 14.6, “Data Dictionary Usage Differences”. Also, for upgrades to MySQL 8.0, the upgrade procedure differs somewhat from previous MySQL releases and requires that you verify the upgrade readiness of your installation by checking specific prerequisites. For more information, see Section 2.10.1, “Upgrading MySQL”, particularly Verifying Upgrade Prerequisites for Your MySQL 5.7 Installation.

14.1 Data Dictionary Schema

Data dictionary tables are invisible and cannot be accessed directly. However, MySQL supports access to data stored in data dictionary tables through `INFORMATION_SCHEMA` tables and `SHOW` statements. For an overview of the tables that comprise the data dictionary, see Data Dictionary Tables.

MySQL system tables still exist in MySQL 8.0 and can be viewed by issuing a `SHOW TABLES` statement on the `mysql` system database. Generally, the difference between MySQL system tables and data dictionary tables is that system tables contain auxiliary data such as time zone and help

information, whereas data dictionary tables contain data required to execute SQL queries. MySQL system tables and data dictionary tables also differ in how they are upgraded. Upgrading MySQL system tables requires running `mysql_upgrade`. Data dictionary upgrades are managed by the MySQL server.

14.2 Removal of File-based Metadata Storage

In previous MySQL releases, dictionary data was partially stored in metadata files. Issues with file-based metadata storage included expensive file scans, susceptibility to file system-related bugs, complex code for handling of replication and crash recovery failure states, and a lack of extensibility that made it difficult to add metadata for new features and relational objects.

The metadata files listed below are removed from MySQL. Unless otherwise noted, data previously stored in metadata files is now stored in data dictionary tables.

- `.frm` files: Table metadata files. With the removal of `.frm` files:
  - The 64KB table definition size limit imposed by the `.frm` file structure is removed.
  - the `INFORMATION_SCHEMA.TABLES VERSION` column now reports a hardcoded value of `10`, which is the last `frm` file version used in MySQL 5.7.
- `.par` files: Partition definition files. `.par` files were removed in MySQL 5.7.6 with the introduction of native partitioning for `InnoDB` tables.
- `.TRN` files: Trigger namespace files.
- `.TRG` files: Trigger parameter files.
- `.isl` files: `InnoDB` Symbolic Link files containing the location of file-per-table tablespace files created outside of the MySQL data directory.
- `db.opt` files: Database configuration files. These files, one per database directory, contained database default character set attributes.

14.3 Transactional Storage of Dictionary Data

The data dictionary schema stores dictionary data in transactional (`InnoDB`) tables. Data dictionary tables are located in the `mysql` database together with non-data dictionary system tables.

As of MySQL 8.0.2, data dictionary tables are created in a single `InnoDB` tablespace named `mysql.ibd` in the MySQL data directory. The `mysql.ibd` tablespace file must reside in the MySQL data directory and its name cannot be modified or used by another tablespace. Previously, these tables were created in individual tablespace files in the `mysql` database directory.

Dictionary data is protected by the same commit, rollback, and crash-recovery capabilities that protect user data stored in `InnoDB` tables.

14.4 Dictionary Object Cache

The dictionary object cache is a shared global cache that stores previously accessed data dictionary objects in memory to enable object reuse and minimize disk I/O. Similar to other cache mechanisms used by MySQL, the dictionary object cache uses an LRU-based eviction strategy to evict least recently used objects from memory.

The dictionary object cache comprises cache partitions that store different object types. Some cache partition size limits are configurable, whereas others are hardcoded.

- **tablespace definition cache partition**: Stores tablespace definition objects. The `tablespace_definition_cache` option sets a limit for the number of tablespace definition objects that can be stored in the dictionary object cache. The default value is 256.

- **schema definition cache partition**: Stores schema definition objects. The `schema_definition_cache` option sets a limit for the number of schema definition objects that can be stored in the dictionary object cache. The default value is 256.
- **table definition cache partition**: Stores table definition objects. The object limit is set to the value of `max_connections`, which has a default value of 151.

  The table definition cache partition exists in parallel with the table definition cache that is configured using the `table_definition_cache` configuration option. Both caches store table definitions but serve different parts of the MySQL server. Objects in one cache have no dependence on the existence of objects in the other.
- **stored program definition cache partition**: Stores stored program definition objects. The `stored_program_definition_cache` option sets a limit for the number of stored program definition objects that can be stored in the dictionary object cache. The default value is 256.

  The stored program definition cache partition exists in parallel with the stored procedure and stored function caches that are configured using the `stored_program_cache` option.

  The `stored_program_cache` option sets a soft upper limit for the number of cached stored procedures or functions per connection, and the limit is checked each time a connection executes a stored procedure or function. The stored program definition cache partition, on the other hand, is a shared cache that stores stored program definition objects for other purposes. The existence of objects in the stored program definition cache partition has no dependence on the existence of objects in the stored procedure cache or stored function cache, and vice versa.
- **character set definition cache partition**: Stores character set definition objects and has a hardcoded object limit of 256.
- **collation definition cache partition**: Stores collation definition objects and has a hardcoded object limit of 256.

For information about valid values for dictionary object cache configuration options, refer to Section 5.1.5, “Server System Variables”.

14.5 INFORMATION_SCHEMA and Data Dictionary Integration

With the introduction of the data dictionary, the following `INFORMATION_SCHEMA` tables are implemented as views on data dictionary tables:

- `CHARACTER_SETS`
- `COLLATIONS`
- `COLLATION_CHARACTER_SET_APPLICABILITY`
- `COLUMNS`
- `EVENTS`
- `KEY_COLUMN_USAGE`
- `PARAMETERS`
- `ROUTINES`
- `SCHEMATA`
- `STATISTICS`
- `ST_SPATIAL_REFERENCE_SYSTEMS`
- `TABLES`

- TABLE_CONSTRAINTS
- TRIGGERS
- VIEWS

Queries on those tables are now more efficient because they obtain information from data dictionary tables rather than by other, slower means. In particular, for each INFORMATION_SCHEMA table that is a view on data dictionary tables:

- The server no longer must create a temporary table for each query of the INFORMATION_SCHEMA table.
- When the underlying data dictionary tables store values previously obtained by directory scans (for example, to enumerate database names or table names within databases) or file-opening operations (for example, to read information from .frm files), INFORMATION_SCHEMA queries for those values now use table lookups instead. (Additionally, even for a non-view INFORMATION_SCHEMA table, values such as database and table names are retrieved by lookups from the data dictionary and do not require directory or file scans.)
- Indexes on the underlying data dictionary tables permit the optimizer to construct efficient query execution plans, something not true for the previous implementation that processed the INFORMATION_SCHEMA table using a temporary table per query.

The preceding improvements also apply to SHOW statements that display information corresponding to the INFORMATION_SCHEMA tables that are views on data dictionary tables. For example, SHOW DATABASES displays the same information as the SCHEMATA table.

In addition to the introduction of views on data dictionary tables, table metadata contained in the STATISTICS and TABLES tables is now cached to improve INFORMATION_SCHEMA query performance. Caching of table metadata is controlled by the information_schema_stats configuration option, which is set to CACHED by default. Cached table metadata is updated by issuing an ANALYZE TABLE statement.

information_schema_stats can be set to LATEST to have INFORMATION_SCHEMA queries retrieve the latest metadata directly from the storage engine, which is not as fast as retrieving cached table metadata.

For more information, see Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”.

14.6 Data Dictionary Usage Differences

Use of a data dictionary-enabled MySQL server entails some operational differences compared to a server that does not have a data dictionary:

- Previously, enabling the innodb_read_only system variable prevented creating and dropping tables only for the InnoDB storage. As of MySQL 8.0, enabling innodb_read_only prevents these operations for all storage engines. Table creation and drop operations for any storage engine modify data dictionary tables in the mysql system database, but those tables use the InnoDB storage engine and cannot be modified when innodb_read_only is enabled. The same principle applies to other table operations that require modifying data dictionary tables. Examples:
  - ANALYZE TABLE fails because it updates table statistics, which are stored in the data dictionary.
  - ALTER TABLE *tbl_name* ENGINE=*engine_name* fails because it updates the storage engine designation, which is stored in the data dictionary.

Note

Enabling innodb_read_only also has important implications for non-data dictionary tables in the mysql system database. For details, see the

description of `innodb_read_only` in Section 15.13, “InnoDB Startup Options and System Variables”

- Previously, tables in the `mysql` system database were visible to DML and DDL statements. As of MySQL 8.0, data dictionary tables are invisible and cannot be modified or queried directly. However, in most cases there are corresponding `INFORMATION_SCHEMA` tables that can be queried instead. This enables the underlying data dictionary tables to be changed as server development proceeds, while maintaining a stable `INFORMATION_SCHEMA` interface for application use.

- `INFORMATION_SCHEMA` tables in MySQL 8.0 are closely tied to the data dictionary, resulting in several usage differences:

  - Previously, `INFORMATION_SCHEMA` queries for table statistics in the `STATISTICS` and `TABLES` tables retrieved statistics directly from storage engines. As of MySQL 8.0, the default is to use cached table statistics, which is more efficient. However, when the server starts, the cached statistics are `NULL` and are not updated for a given table until `ANALYZE TABLE` is run. Also, the statistics go out of date to the extent that the table is modified thereafter. To use the latest table statistics obtained from storage engines (the pre-8.0 behavior), change the `information_schema_stats` system variable from `CACHED` to `LATEST`. For more information, see Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”.

  - Several `INFORMATION_SCHEMA` tables are views on data dictionary tables, which enables the optimizer to use indexes on those underlying tables. Consequently, depending on optimizer choices, the row order of results for `INFORMATION_SCHEMA` queries might differ from previous results. If a query result must have specific row ordering characteristics, include an `ORDER BY` clause.

  - `mysqldump` and `mysqlpump` no longer dump the `INFORMATION_SCHEMA` database, even if explicitly named on the command line.

  - `CREATE TABLE dst_tbl LIKE src_tbl` requires that `src_tbl` be a base table and fails if it is an `INFORMATION_SCHEMA` table that is a view on data dictionary tables.

  - Previously, result set headers of columns selected from `INFORMATION_SCHEMA` tables used the capitalization specified in the query. This query produces a result set with a header of `table_name`:

    ```
    SELECT table_name FROM INFORMATION_SCHEMA.TABLES;
    ```

    As of MySQL 8.0, these headers are capitalized; the preceding query produces a result set with a header of `TABLE_NAME`. If necessary, a column alias can be used to achieve a different lettercase. For example:

    ```
    SELECT table_name AS 'table_name' FROM INFORMATION_SCHEMA.TABLES;
    ```

- The data directory affects how `mysqldump` and `mysqlpump` dump information from the `mysql` system database:

  - Previously, it was possible to dump all tables in the `mysql` system database. As of MySQL 8.0, `mysqldump` and `mysqlpump` dump only non-data dictionary tables in that database.

  - Previously, the `--routines` and `--events` options were not required to include stored routines and events when using the `--all-databases` option: The dump included the `mysql` system database, and therefore also the `proc` and `event` tables containing stored routine and event definitions. As of MySQL 8.0, the `event` and `proc` tables are not used. Definitions for the corresponding objects are stored in data dictionary tables, but those tables are not dumped. To include stored routines and events in a dump made using `--all-databases`, use the `--routines` and `--events` options explicitly.

  - Previously, the `--routines` option required the `SELECT` privilege for the `proc` table. As of MySQL 8.0, that table is not used; `--routines` requires the global `SELECT` privilege instead.
  - Previously, it was possible to dump stored routine and event definitions together with their creation and modification timestamps, by dumping the `proc` and `event` tables. As of MySQL 8.0, those tables are not used, so it is not possible to dump timestamps.
- Previously, creating a stored routine that contains illegal characters produced a warning. As of MySQL 8.0, this is an error.

14.7 Data Dictionary Limitations

This section describes temporary limitations introduced with the MySQL data dictionary.

- Manual creation of database directories under the data directory (for example, with `mkdir`) is unsupported. Manually created database directories are not recognized by the MySQL Server.
- Moving data stored in `MyISAM` tables by copying and moving `MyISAM` data files is unsupported. Tables moved using this method are not discovered by the server.
- Simple backup and restore of individual `MyISAM` tables using copied data files is unsupported.
- `TRUNCATE TABLE`, which is mapped to `DROP TABLE` and `CREATE TABLE` in MySQL 8.0, is temporarily non-atomic. A server exit during a `TRUNCATE TABLE` operation can result in a dropped table. It may also result in orphaned foreign key entries in the `InnoDB` `SYS_FOREIGN` and `SYS_FOREIGN_COLS` dictionary tables if the table contains foreign key constraints.
- DDL operations take longer due to writing to storage, undo logs, and redo logs instead of `.frm` files.
- The period of time during which DDL operations are vulnerable to a server exit is longer due to writing to storage, undo logs, and redo logs instead of `.frm` files. The minimal risk introduced by this limitation is mostly applicable to restore operations and other operations that load numerous new tables.
- With the removal of `.isl` files in MySQL 8.0, offline relocation of file-per-table tablespaces created outside of the MySQL data directory is not supported.
- If MySQL exits in the middle of an `ALTER TABLE` operation on an `InnoDB` table, you may be left with an orphan intermediate or temporary table. Removing orphan intermediate and temporary tables is unsupported.

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15.1 Introduction to InnoDB

`InnoDB` is a general-purpose storage engine that balances high reliability and high performance. In MySQL 8.0, `InnoDB` is the default MySQL storage engine. Unless you have configured a different default storage engine, issuing a `CREATE TABLE` statement without an `ENGINE=` clause creates an `InnoDB` table.

Key Advantages of InnoDB

Key advantages of `InnoDB` include:

- Its DML operations follow the ACID model, with transactions featuring commit, rollback, and crash-recovery capabilities to protect user data. See Section 15.2, “InnoDB and the ACID Model” for more information.
- Row-level locking and Oracle-style consistent reads increase multi-user concurrency and performance. See Section 15.5, “InnoDB Locking and Transaction Model” for more information.
- `InnoDB` tables arrange your data on disk to optimize queries based on primary keys. Each `InnoDB` table has a primary key index called the clustered index that organizes the data to minimize I/O for primary key lookups. See Section 15.8.2.1, “Clustered and Secondary Indexes” for more information.
- To maintain data integrity, `InnoDB` supports `FOREIGN KEY` constraints. With foreign keys, inserts, updates, and deletes are checked to ensure they do not result in inconsistencies across different tables. See Section 15.8.1.6, “InnoDB and FOREIGN KEY Constraints” for more information.

Table 15.1 InnoDB Storage Engine Features

Storage limits	64TB	**Transactions**	Yes	**Locking granularity**	Row
MVCC	Yes	**Geospatial data type support**	Yes	**Geospatial indexing support**	Yes[a]
B-tree indexes	Yes	**T-tree indexes**	No	**Hash indexes**	No[b]
Full-text search indexes	Yes[c]	**Clustered indexes**	Yes	**Data caches**	Yes
Index caches	Yes	**Compressed data**	Yes[d]	**Encrypted data**[e]	Yes
Cluster database support	No	**Replication support**[f]	Yes	**Foreign key support**	Yes
Backup / point-in-time recovery[g]	Yes	**Query cache support**	Yes	**Update statistics for data dictionary**	Yes

[a] InnoDB support for geospatial indexing is available in MySQL 5.7.5 and later.
[b] InnoDB utilizes hash indexes internally for its Adaptive Hash Index feature.
[c] InnoDB support for FULLTEXT indexes is available in MySQL 5.6.4 and later.
[d] Compressed InnoDB tables require the InnoDB Barracuda file format.
[e] Implemented in the server (via encryption functions). Data-at-rest tablespace encryption is available in MySQL 5.7 and later.
[f] Implemented in the server, rather than in the storage engine.
[g] Implemented in the server, rather than in the storage engine.

To compare the features of `InnoDB` with other storage engines provided with MySQL, see the *Storage Engine Features* table in Chapter 16, *Alternative Storage Engines*.

InnoDB Enhancements and New Features

For information about `InnoDB` enhancements and new features in MySQL 8.0, refer to:

- The InnoDB enhancements list in Section 1.4, “What Is New in MySQL 8.0”.
- The Release Notes.

Additional InnoDB Information and Resources

- For InnoDB-related terms and definitions, see MySQL Glossary.
- For a forum dedicated to the InnoDB storage engine, see MySQL Forums::InnoDB.
- InnoDB is published under the same GNU GPL License Version 2 (of June 1991) as MySQL. For more information on MySQL licensing, see http://www.mysql.com/company/legal/licensing/.

15.1.1 Benefits of Using InnoDB Tables

If you use MyISAM tables but are not committed to them for technical reasons, you may find InnoDB tables beneficial for the following reasons:

- If your server crashes because of a hardware or software issue, regardless of what was happening in the database at the time, you don't need to do anything special after restarting the database. InnoDB crash recovery automatically finalizes any changes that were committed before the time of the crash, and undoes any changes that were in process but not committed. Just restart and continue where you left off.
- The InnoDB storage engine maintains its own buffer pool that caches table and index data in main memory as data is accessed. Frequently used data is processed directly from memory. This cache applies to many types of information and speeds up processing. On dedicated database servers, up to 80% of physical memory is often assigned to the InnoDB buffer pool.
- If you split up related data into different tables, you can set up foreign keys that enforce referential integrity. Update or delete data, and the related data in other tables is updated or deleted automatically. Try to insert data into a secondary table without corresponding data in the primary table, and the bad data gets kicked out automatically.
- If data becomes corrupted on disk or in memory, a checksum mechanism alerts you to the bogus data before you use it.
- When you design your database with appropriate primary key columns for each table, operations involving those columns are automatically optimized. It is very fast to reference the primary key columns in WHERE clauses, ORDER BY clauses, GROUP BY clauses, and join operations.
- Inserts, updates, and deletes are optimized by an automatic mechanism called change buffering. InnoDB not only allows concurrent read and write access to the same table, it caches changed data to streamline disk I/O.
- Performance benefits are not limited to giant tables with long-running queries. When the same rows are accessed over and over from a table, a feature called the Adaptive Hash Index takes over to make these lookups even faster, as if they came out of a hash table.
- You can compress tables and associated indexes.
- You can create and drop indexes with much less impact on performance and availability.
- Truncating a file-per-table tablespace is very fast, and can free up disk space for the operating system to reuse, rather than freeing up space within the system tablespace that only InnoDB could reuse.
- The storage layout for table data is more efficient for BLOB and long text fields, with the DYNAMIC row format.
- You can monitor the internal workings of the storage engine by querying INFORMATION_SCHEMA tables.

- You can monitor the performance details of the storage engine by querying Performance Schema tables.
- You can freely mix `InnoDB` tables with tables from other MySQL storage engines, even within the same statement. For example, you can use a join operation to combine data from `InnoDB` and `MEMORY` tables in a single query.
- `InnoDB` has been designed for CPU efficiency and maximum performance when processing large data volumes.
- `InnoDB` tables can handle large quantities of data, even on operating systems where file size is limited to 2GB.

For `InnoDB`-specific tuning techniques you can apply in your application code, see Section 8.5, “Optimizing for InnoDB Tables”.

15.1.2 Best Practices for InnoDB Tables

This section describes best practices when using `InnoDB` tables.

- Specifying a primary key for every table using the most frequently queried column or columns, or an auto-increment value if there is no obvious primary key.
- Using joins wherever data is pulled from multiple tables based on identical ID values from those tables. For fast join performance, define foreign keys on the join columns, and declare those columns with the same data type in each table. Adding foreign keys ensures that referenced columns are indexed, which can improve performance. Foreign keys also propagate deletes or updates to all affected tables, and prevent insertion of data in a child table if the corresponding IDs are not present in the parent table.
- Turning off autocommit. Committing hundreds of times a second puts a cap on performance (limited by the write speed of your storage device).
- Grouping sets of related DML operations into transactions, by bracketing them with `START TRANSACTION` and `COMMIT` statements. While you don't want to commit too often, you also don't want to issue huge batches of `INSERT`, `UPDATE`, or `DELETE` statements that run for hours without committing.
- Not using `LOCK TABLES` statements. `InnoDB` can handle multiple sessions all reading and writing to the same table at once, without sacrificing reliability or high performance. To get exclusive write access to a set of rows, use the `SELECT ... FOR UPDATE` syntax to lock just the rows you intend to update.
- Enabling the `innodb_file_per_table` option to put the data and indexes for individual tables into separate files, instead of in a single giant system tablespace. This setting is required to use some of the other features, such as table compression and fast truncation.

  The `innodb_file_per_table` option is enabled by default as of MySQL 5.6.6.
- Evaluating whether your data and access patterns benefit from the `InnoDB` table compression feature (`ROW_FORMAT=COMPRESSED`) on the `CREATE TABLE` statement. You can compress `InnoDB` tables without sacrificing read/write capability.
- Running your server with the option `--sql_mode=NO_ENGINE_SUBSTITUTION` to prevent tables being created with a different storage engine if there is an issue with the engine specified in the `ENGINE=` clause of `CREATE TABLE`.

15.1.3 Verifying that InnoDB is the Default Storage Engine

To verify that `InnoDB` is the default storage engine, issue the `SHOW ENGINES` command to view the different MySQL storage engines. Look for `DEFAULT` in the `InnoDB` line. Alternatively, query the `INFORMATION_SCHEMA ENGINES` table.

15.1.4 Testing and Benchmarking with InnoDB

If `InnoDB` is not your default storage engine, you can determine if your database server or applications work correctly with `InnoDB` by restarting the server with `--default-storage-engine=InnoDB` defined on the command line or with `default-storage-engine=innodb` defined in the `[mysqld]` section of the `my.cnf` configuration file.

Since changing the default storage engine only affects new tables as they are created, run all your application installation and setup steps to confirm that everything installs properly. Then exercise all the application features to make sure all the data loading, editing, and querying features work. If a table relies on some `MyISAM`-specific feature, you'll receive an error; add the `ENGINE=MyISAM` clause to the `CREATE TABLE` statement to avoid the error.

If you did not make a deliberate decision about the storage engine, and you just want to preview how certain tables work when they're created under `InnoDB`, issue the command `ALTER TABLE table_name ENGINE=InnoDB;` for each table. Or, to run test queries and other statements without disturbing the original table, make a copy like so:

```
CREATE TABLE InnoDB_Table (...) ENGINE=InnoDB AS SELECT * FROM MyISAM_Table;
```

To get a true idea of the performance with a full application under a realistic workload, install the latest MySQL server and run benchmarks.

Test the full application lifecycle, from installation, through heavy usage, and server restart. Kill the server process while the database is busy to simulate a power failure, and verify that the data is recovered successfully when you restart the server.

Test any replication configurations, especially if you use different MySQL versions and options on the master and the slaves.

15.2 InnoDB and the ACID Model

The ACID model is a set of database design principles that emphasize aspects of reliability that are important for business data and mission-critical applications. MySQL includes components such as the `InnoDB` storage engine that adhere closely to the ACID model, so that data is not corrupted and results are not distorted by exceptional conditions such as software crashes and hardware malfunctions. When you rely on ACID-compliant features, you do not need to reinvent the wheel of consistency checking and crash recovery mechanisms. In cases where you have additional software safeguards, ultra-reliable hardware, or an application that can tolerate a small amount of data loss or inconsistency, you can adjust MySQL settings to trade some of the ACID reliability for greater performance or throughput.

The following sections discuss how MySQL features, in particular the `InnoDB` storage engine, interact with the categories of the ACID model:

- **A**: atomicity.
- **C**: consistency.
- **I:**: isolation.
- **D**: durability.

Atomicity

The **atomicity** aspect of the ACID model mainly involves `InnoDB` transactions. Related MySQL features include:

- Autocommit setting.
- `COMMIT` statement.
- `ROLLBACK` statement.
- Operational data from the `INFORMATION_SCHEMA` tables.

Consistency

The **consistency** aspect of the ACID model mainly involves internal `InnoDB` processing to protect data from crashes. Related MySQL features include:

- `InnoDB` doublewrite buffer.
- `InnoDB` crash recovery.

Isolation

The **isolation** aspect of the ACID model mainly involves `InnoDB` transactions, in particular the isolation level that applies to each transaction. Related MySQL features include:

- Autocommit setting.
- `SET ISOLATION LEVEL` statement.
- The low-level details of `InnoDB` locking. During performance tuning, you see these details through `INFORMATION_SCHEMA` tables.

Durability

The **durability** aspect of the ACID model involves MySQL software features interacting with your particular hardware configuration. Because of the many possibilities depending on the capabilities of your CPU, network, and storage devices, this aspect is the most complicated to provide concrete guidelines for. (And those guidelines might take the form of buy “new hardware”.) Related MySQL features include:

- `InnoDB` doublewrite buffer, turned on and off by the `innodb_doublewrite` configuration option.
- Configuration option `innodb_flush_log_at_trx_commit`.
- Configuration option `sync_binlog`.
- Configuration option `innodb_file_per_table`.
- Write buffer in a storage device, such as a disk drive, SSD, or RAID array.
- Battery-backed cache in a storage device.
- The operating system used to run MySQL, in particular its support for the `fsync()` system call.
- Uninterruptible power supply (UPS) protecting the electrical power to all computer servers and storage devices that run MySQL servers and store MySQL data.
- Your backup strategy, such as frequency and types of backups, and backup retention periods.
- For distributed or hosted data applications, the particular characteristics of the data centers where the hardware for the MySQL servers is located, and network connections between the data centers.

15.3 InnoDB Multi-Versioning

InnoDB is a multi-versioned storage engine: it keeps information about old versions of changed rows, to support transactional features such as concurrency and rollback. This information is stored in the tablespace in a data structure called a rollback segment (after an analogous data structure in Oracle). InnoDB uses the information in the rollback segment to perform the undo operations needed in a transaction rollback. It also uses the information to build earlier versions of a row for a consistent read.

Internally, InnoDB adds three fields to each row stored in the database. A 6-byte DB_TRX_ID field indicates the transaction identifier for the last transaction that inserted or updated the row. Also, a deletion is treated internally as an update where a special bit in the row is set to mark it as deleted. Each row also contains a 7-byte DB_ROLL_PTR field called the roll pointer. The roll pointer points to an undo log record written to the rollback segment. If the row was updated, the undo log record contains the information necessary to rebuild the content of the row before it was updated. A 6-byte DB_ROW_ID field contains a row ID that increases monotonically as new rows are inserted. If InnoDB generates a clustered index automatically, the index contains row ID values. Otherwise, the DB_ROW_ID column does not appear in any index.

Undo logs in the rollback segment are divided into insert and update undo logs. Insert undo logs are needed only in transaction rollback and can be discarded as soon as the transaction commits. Update undo logs are used also in consistent reads, but they can be discarded only after there is no transaction present for which InnoDB has assigned a snapshot that in a consistent read could need the information in the update undo log to build an earlier version of a database row.

Commit your transactions regularly, including those transactions that issue only consistent reads. Otherwise, InnoDB cannot discard data from the update undo logs, and the rollback segment may grow too big, filling up your tablespace.

The physical size of an undo log record in the rollback segment is typically smaller than the corresponding inserted or updated row. You can use this information to calculate the space needed for your rollback segment.

In the InnoDB multi-versioning scheme, a row is not physically removed from the database immediately when you delete it with an SQL statement. InnoDB only physically removes the corresponding row and its index records when it discards the update undo log record written for the deletion. This removal operation is called a purge, and it is quite fast, usually taking the same order of time as the SQL statement that did the deletion.

If you insert and delete rows in smallish batches at about the same rate in the table, the purge thread can start to lag behind and the table can grow bigger and bigger because of all the “dead” rows, making everything disk-bound and very slow. In such a case, throttle new row operations, and allocate more resources to the purge thread by tuning the innodb_max_purge_lag system variable. See Section 15.13, “InnoDB Startup Options and System Variables” for more information.

Multi-Versioning and Secondary Indexes

InnoDB multiversion concurrency control (MVCC) treats secondary indexes differently than clustered indexes. Records in a clustered index are updated in-place, and their hidden system columns point undo log entries from which earlier versions of records can be reconstructed. Unlike clustered index records, secondary index records do not contain hidden system columns nor are they updated in-place.

When a secondary index column is updated, old secondary index records are delete-marked, new records are inserted, and delete-marked records are eventually purged. When a secondary index record is delete-marked or the secondary index page is updated by a newer transaction, InnoDB looks up the database record in the clustered index. In the clustered index, the record's DB_TRX_ID is checked, and the correct version of the record is retrieved from the undo log if the record was modified after the reading transaction was initiated.

If a secondary index record is marked for deletion or the secondary index page is updated by a newer transaction, the covering index technique is not used. Instead of returning values from the index structure, InnoDB looks up the record in the clustered index.

However, if the index condition pushdown (ICP) optimization is enabled, and parts of the `WHERE` condition can be evaluated using only fields from the index, the MySQL server still pushes this part of the `WHERE` condition down to the storage engine where it is evaluated using the index. If no matching records are found, the clustered index lookup is avoided. If matching records are found, even among delete-marked records, `InnoDB` looks up the record in the clustered index.

15.4 InnoDB Architecture

This section provides an introduction to the major components of the `InnoDB` storage engine architecture.

15.4.1 Buffer Pool

The buffer pool is an area in main memory where `InnoDB` caches table and index data as data is accessed. The buffer pool allows frequently used data to be processed directly from memory, which speeds up processing. On dedicated database servers, up to 80% of physical memory is often assigned to the `InnoDB` buffer pool.

For efficiency of high-volume read operations, the buffer pool is divided into pages that can potentially hold multiple rows. For efficiency of cache management, the buffer pool is implemented as a linked list of pages; data that is rarely used is aged out of the cache, using a variation of the LRU algorithm.

For more information, see Section 15.6.3.1, “The InnoDB Buffer Pool”, and Section 15.6.3, “InnoDB Buffer Pool Configuration”.

15.4.2 Change Buffer

The change buffer is a special data structure that caches changes to secondary index pages when affected pages are not in the buffer pool. The buffered changes, which may result from `INSERT`, `UPDATE`, or `DELETE` operations (DML), are merged later when the pages are loaded into the buffer pool by other read operations.

Unlike clustered indexes, secondary indexes are usually non-unique, and inserts into secondary indexes happen in a relatively random order. Similarly, deletes and updates may affect secondary index pages that are not adjacently located in an index tree. Merging cached changes at a later time, when affected pages are read into the buffer pool by other operations, avoids substantial random access I/O that would be required to read-in secondary index pages from disk.

Periodically, the purge operation that runs when the system is mostly idle, or during a slow shutdown, writes the updated index pages to disk. The purge operation can write disk blocks for a series of index values more efficiently than if each value were written to disk immediately.

Change buffer merging may take several hours when there are numerous secondary indexes to update and many affected rows. During this time, disk I/O is increased, which can cause a significant slowdown for disk-bound queries. Change buffer merging may also continue to occur after a transaction is committed. In fact, change buffer merging may continue to occur after a server shutdown and restart (see Section 15.20.2, “Forcing InnoDB Recovery” for more information).

In memory, the change buffer occupies part of the `InnoDB` buffer pool. On disk, the change buffer is part of the system tablespace, so that index changes remain buffered across database restarts.

The type of data cached in the change buffer is governed by the `innodb_change_buffering` configuration option. For more information, see Section 15.6.4, “Configuring InnoDB Change Buffering”. You can also configure the maximum change buffer size. For more information, see Section 15.6.4.1, “Configuring the Change Buffer Maximum Size”.

Change buffering is not supported for a secondary index if the index contains a descending index column or if the primary key includes a descending index column.

Monitoring the Change Buffer

The following options are available for change buffer monitoring:

- `InnoDB` Standard Monitor output includes status information for the change buffer. To view monitor data, issue the `SHOW ENGINE INNODB STATUS` command.

```
mysql> SHOW ENGINE INNODB STATUS\G
```

Change buffer status information is located under the `INSERT BUFFER AND ADAPTIVE HASH INDEX` heading and appears similar to the following:

```
-------------------------------------
INSERT BUFFER AND ADAPTIVE HASH INDEX
-------------------------------------
Ibuf: size 1, free list len 0, seg size 2, 0 merges
merged operations:
 insert 0, delete mark 0, delete 0
discarded operations:
 insert 0, delete mark 0, delete 0
Hash table size 4425293, used cells 32, node heap has 1 buffer(s)
13577.57 hash searches/s, 202.47 non-hash searches/s
```

For more information, see Section 15.16.3, “InnoDB Standard Monitor and Lock Monitor Output”.

- The `INFORMATION_SCHEMA.INNODB_METRICS` table provides most of the data points found in `InnoDB` Standard Monitor output, plus other data points. To view change buffer metrics and a description of each, issue the following query:

```
mysql> SELECT NAME, COMMENT FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME LIKE '%ibuf%'\G
```

For `INNODB_METRICS` table usage information, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

- The `INFORMATION_SCHEMA.INNODB_BUFFER_PAGE` table provides metadata about each page in the buffer pool, including change buffer index and change buffer bitmap pages. Change buffer pages are identified by `PAGE_TYPE`. `IBUF_INDEX` is the page type for change buffer index pages, and `IBUF_BITMAP` is the page type for change buffer bitmap pages.

Warning

Querying the `INNODB_BUFFER_PAGE` table can introduce significant performance overhead. To avoid impacting performance, reproduce the issue you want to investigate on a test instance and run your queries on the test instance.

For example, you can query the `INNODB_BUFFER_PAGE` table to determine the approximate number of `IBUF_INDEX` and `IBUF_BITMAP` pages as a percentage of total buffer pool pages.

```
SELECT
(SELECT COUNT(*) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE PAGE_TYPE LIKE 'IBUF%'
) AS change_buffer_pages,
(
SELECT COUNT(*)
FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
) AS total_pages,
(
SELECT ((change_buffer_pages/total_pages)*100)
) AS change_buffer_page_percentage;
+---------------------+-------------+-------------------------------+
| change_buffer_pages | total_pages | change_buffer_page_percentage |
+---------------------+-------------+-------------------------------+
|                  25 |        8192 |                        0.3052 |
+---------------------+-------------+-------------------------------+
```

For information about other data provided by the `INNODB_BUFFER_PAGE` table, see Section 24.31.1, "The INFORMATION_SCHEMA INNODB_BUFFER_PAGE Table". For related usage information, see Section 15.14.5, "InnoDB INFORMATION_SCHEMA Buffer Pool Tables".

- Performance Schema provides change buffer mutex wait instrumentation for advanced performance monitoring. To view change buffer instrumentation, issue the following query:

```
mysql> SELECT * FROM performance_schema.setup_instruments
WHERE NAME LIKE '%wait/synch/mutex/innodb/ibuf%';
+-----------------------------------------------------+---------+-------+
| NAME                                                | ENABLED | TIMED |
+-----------------------------------------------------+---------+-------+
| wait/synch/mutex/innodb/ibuf_bitmap_mutex           | YES     | YES   |
| wait/synch/mutex/innodb/ibuf_mutex                  | YES     | YES   |
| wait/synch/mutex/innodb/ibuf_pessimistic_insert_mutex | YES     | YES   |
+-----------------------------------------------------+---------+-------+
```

For information about monitoring `InnoDB` mutex waits, see Section 15.15.2, "Monitoring InnoDB Mutex Waits Using Performance Schema".

15.4.3 Adaptive Hash Index

The adaptive hash index (AHI) lets `InnoDB` perform more like an in-memory database on systems with appropriate combinations of workload and ample memory for the buffer pool, without sacrificing any transactional features or reliability. This feature is enabled by the `innodb_adaptive_hash_index` option, or turned off by `--skip-innodb_adaptive_hash_index` at server startup.

Based on the observed pattern of searches, MySQL builds a hash index using a prefix of the index key. The prefix of the key can be any length, and it may be that only some of the values in the B-tree appear in the hash index. Hash indexes are built on demand for those pages of the index that are often accessed.

If a table fits almost entirely in main memory, a hash index can speed up queries by enabling direct lookup of any element, turning the index value into a sort of pointer. `InnoDB` has a mechanism that monitors index searches. If `InnoDB` notices that queries could benefit from building a hash index, it does so automatically.

With some workloads, the speedup from hash index lookups greatly outweighs the extra work to monitor index lookups and maintain the hash index structure. Sometimes, the read/write lock that guards access to the adaptive hash index can become a source of contention under heavy workloads, such as multiple concurrent joins. Queries with `LIKE` operators and `%` wildcards also tend not to benefit from the AHI. For workloads where the adaptive hash index is not needed, turning it off reduces unnecessary performance overhead. Because it is difficult to predict in advance whether this feature is appropriate for a particular system, consider running benchmarks with it both enabled and disabled, using a realistic workload. The architectural changes in MySQL 5.6 and higher make more workloads suitable for disabling the adaptive hash index than in earlier releases, although it is still enabled by default.

The adaptive hash index search system is partitioned. Each index is bound to a specific partition, and each partition is protected by a separate latch. Partitioning is controlled by the `innodb_adaptive_hash_index_parts` configuration option. The `innodb_adaptive_hash_index_parts` option is set to 8 by default. The maximum setting is 512.

The hash index is always built based on an existing B-tree index on the table. `InnoDB` can build a hash index on a prefix of any length of the key defined for the B-tree, depending on the pattern of searches that `InnoDB` observes for the B-tree index. A hash index can be partial, covering only those pages of the index that are often accessed.

You can monitor the use of the adaptive hash index and the contention for its use in the `SEMAPHORES` section of the output of the `SHOW ENGINE INNODB STATUS` command. If you see many threads

waiting on an RW-latch created in `btr0sea.c`, then it might be useful to disable adaptive hash indexing.

For more information about the performance characteristics of hash indexes, see Section 8.3.8, “Comparison of B-Tree and Hash Indexes”.

15.4.4 Redo Log Buffer

The redo log buffer is the memory area that holds data to be written to the redo log. Redo log buffer size is defined by the `innodb_log_buffer_size` configuration option. The redo log buffer is periodically flushed to the log file on disk. A large redo log buffer enables large transactions to run without the need to write redo log to disk before the transactions commit. Thus, if you have transactions that update, insert, or delete many rows, making the log buffer larger saves disk I/O.

The `innodb_flush_log_at_trx_commit` option controls how the contents of the redo log buffer are written to the log file. The `innodb_flush_log_at_timeout` option controls redo log flushing frequency.

15.4.5 System Tablespace

The `InnoDB` system tablespace contains the `InnoDB` data dictionary (metadata for `InnoDB`-related objects) and is the storage area for the doublewrite buffer, the change buffer, and undo logs. The system tablespace also contains table and index data for any user-created tables that are created in the system tablespace. The system tablespace is considered a shared tablespace since it is shared by multiple tables.

The system tablespace is represented by one or more data files. By default, one system data file, named `ibdata1`, is created in the MySQL `data` directory. The size and number of system data files is controlled by the `innodb_data_file_path` startup option.

For related information, see Section 15.6.1, “InnoDB Startup Configuration”, and Section 15.7.1, “Resizing the InnoDB System Tablespace”.

15.4.6 Doublewrite Buffer

The doublewrite buffer is a storage area located in the system tablespace where `InnoDB` writes pages that are flushed from the `InnoDB` buffer pool, before the pages are written to their proper positions in the data file. Only after flushing and writing pages to the doublewrite buffer, does `InnoDB` write pages to their proper positions. If there is an operating system, storage subsystem, or `mysqld` process crash in the middle of a page write, `InnoDB` can later find a good copy of the page from the doublewrite buffer during crash recovery.

Although data is always written twice, the doublewrite buffer does not require twice as much I/O overhead or twice as many I/O operations. Data is written to the doublewrite buffer itself as a large sequential chunk, with a single `fsync()` call to the operating system.

The doublewrite buffer is enabled by default in most cases. To disable the doublewrite buffer, set `innodb_doublewrite` to 0.

If system tablespace files (“ibdata files”) are located on Fusion-io devices that support atomic writes, doublewrite buffering is automatically disabled and Fusion-io atomic writes are used for all data files. Because the doublewrite buffer setting is global, doublewrite buffering is also disabled for data files residing on non-Fusion-io hardware. This feature is only supported on Fusion-io hardware and is only enabled for Fusion-io NVMFS on Linux. To take full advantage of this feature, an `innodb_flush_method` setting of `O_DIRECT` is recommended.

15.4.7 Undo Logs

An undo log is a collection of undo log records associated with a single transaction. An undo log record contains information about how to undo the latest change by a transaction to a clustered index

record. If another transaction needs to see the original data (as part of a consistent read operation), the unmodified data is retrieved from the undo log records. Undo logs exist within undo log segments, which are contained within rollback segments. Rollback segments reside in the system tablespace, temporary tablespace, and undo tablespaces. For more information about undo tablespaces, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”. For information about multi-versioning, see Section 15.3, “InnoDB Multi-Versioning”.

Prior to MySQL 8.0.2, `InnoDB` supports a maximum of 128 rollback segments. As of MySQL 8.0.2, the system tablespace, temporary tablespace, and each undo tablespace individually support a maximum of 128 rollback segments. The `innodb_rollback_segments` configuration option defines the number of rollback segments. Each rollback segment supports up to 1023 concurrent data-modifying transactions.

15.4.8 File-Per-Table Tablespaces

A file-per-table tablespace is a single-table tablespace that is created in its own data file rather than in the system tablespace. Tables are created in file-per-table tablespaces when the `innodb_file_per_table` option is enabled. Otherwise, `InnoDB` tables are created in the system tablespace. Each file-per-table tablespace is represented by a single `.ibd` data file, which is created in the database directory by default.

File per-table tablespaces support `DYNAMIC` and `COMPRESSED` row formats which support features such as off-page storage for variable length data and table compression. For information about these features, and about other advantages of file-per-table tablespaces, see Section 15.7.4, “InnoDB File-Per-Table Tablespaces”.

15.4.9 General Tablespaces

A shared `InnoDB` tablespace created using `CREATE TABLESPACE` syntax. General tablespaces can be created outside of the MySQL data directory, are capable of holding multiple tables, and support tables of all row formats.

Tables are added to a general tablespace using `CREATE TABLE` *`tbl_name`* `... TABLESPACE [=]` *`tablespace_name`* or `ALTER TABLE` *`tbl_name`* `TABLESPACE [=]` *`tablespace_name`* syntax.

For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

15.4.10 Undo Tablespace

An undo tablespace comprises one or more files that contain undo logs. An undo tablespace is created when the `innodb_undo_tablespaces` configuration option is set to a non-zero value. For more information, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

15.4.11 Temporary Tablespace

The temporary tablespace is used for non-compressed `InnoDB` temporary tables and related objects. The `innodb_temp_data_file_path` configuration option defines a relative path for the temporary tablespace data file. If `innodb_temp_data_file_path` is not defined, a single auto-extending 12MB data file named `ibtmp1` is created in the data directory. The temporary tablespace is recreated on each server start and receives a dynamically generated space ID, which helps avoid conflicts with existing space IDs. The temporary tablespace cannot reside on a raw device. Startup is refused if the temporary tablespace cannot be created.

The temporary tablespace is removed on normal shutdown or on an aborted initialization. The temporary tablespace is not removed when a crash occurs. In this case, the database administrator may remove the temporary tablespace manually or restart the server with the same configuration, which removes and recreates the temporary tablespace.

15.4.11.1 Temporary Table Undo Logs

Temporary table undo logs are used for temporary tables and related objects. This type of undo log is not a redo log, as temporary tables are not recovered during crash recovery and do not require redo logs. Temporary table undo logs are, however, used for rollback while the server is running. This special type of non-redo undo log benefits performance by avoiding redo logging I/O for temporary tables and related objects. Temporary table undo logs reside in the temporary tablespace.

Prior to MySQL 8.0.2, 32 rollback segments are assigned for temporary table undo logs for transactions that modify temporary tables and related objects. As of MySQL 8.0.2, the `innodb_rollback_segments` configuration option defines the number of rollback segments used by the temporary tablespace.

15.4.12 Redo Log

The redo log is a disk-based data structure used during crash recovery to correct data written by incomplete transactions. During normal operations, the redo log encodes requests to change `InnoDB` table data that result from SQL statements or low-level API calls. Modifications that did not finish updating the data files before an unexpected shutdown are replayed automatically during initialization, and before the connections are accepted. For information about the role of the redo log in crash recovery, see Section 15.17.2, “InnoDB Recovery”.

By default, the redo log is physically represented on disk as a set of files, named `ib_logfile0` and `ib_logfile1`. MySQL writes to the redo log files in a circular fashion. Data in the redo log is encoded in terms of records affected; this data is collectively referred to as redo. The passage of data through the redo log is represented by an ever-increasing LSN value.

For related information, see:

- Section 15.6.1, “InnoDB Startup Configuration”
- Section 8.5.4, “Optimizing InnoDB Redo Logging”
- Section 15.7.2, “Changing the Number or Size of InnoDB Redo Log Files”
- InnoDB Crash Recovery

15.4.12.1 Group Commit for Redo Log Flushing

`InnoDB`, like any other ACID-compliant database engine, flushes the redo log of a transaction before it is committed. `InnoDB` uses group commit functionality to group multiple such flush requests together to avoid one flush for each commit. With group commit, `InnoDB` issues a single write to the log file to perform the commit action for multiple user transactions that commit at about the same time, significantly improving throughput.

For more information about performance of `COMMIT` and other transactional operations, see Section 8.5.2, “Optimizing InnoDB Transaction Management”.

15.5 InnoDB Locking and Transaction Model

To implement a large-scale, busy, or highly reliable database application, to port substantial code from a different database system, or to tune MySQL performance, it is important to understand `InnoDB` locking and the `InnoDB` transaction model.

This section discusses several topics related to `InnoDB` locking and the `InnoDB` transaction model with which you should be familiar.

- Section 15.5.1, “InnoDB Locking” describes lock types used by `InnoDB`.

- Section 15.5.2, “InnoDB Transaction Model” describes transaction isolation levels and the locking strategies used by each. It also discusses the use of `autocommit`, consistent non-locking reads, and locking reads.
- Section 15.5.3, “Locks Set by Different SQL Statements in InnoDB” discusses specific types of locks set in `InnoDB` for various statements.
- Section 15.5.4, “Phantom Rows” describes how `InnoDB` uses next-key locking to avoid phantom rows.
- Section 15.5.5, “Deadlocks in InnoDB” provides a deadlock example, discusses deadlock detection and rollback, and provides tips for minimizing and handling deadlocks in `InnoDB`.

15.5.1 InnoDB Locking

This section describes lock types used by `InnoDB`.

- Shared and Exclusive Locks
- Intention Locks
- Record Locks
- Gap Locks
- Next-Key Locks
- Insert Intention Locks
- AUTO-INC Locks
- Predicate Locks for Spatial Indexes

Shared and Exclusive Locks

`InnoDB` implements standard row-level locking where there are two types of locks, shared (*S*) locks and exclusive (*X*) locks.

- A shared (*S*) lock permits the transaction that holds the lock to read a row.
- An exclusive (*X*) lock permits the transaction that holds the lock to update or delete a row.

If transaction `T1` holds a shared (*S*) lock on row `r`, then requests from some distinct transaction `T2` for a lock on row `r` are handled as follows:

- A request by `T2` for an *S* lock can be granted immediately. As a result, both `T1` and `T2` hold an *S* lock on `r`.
- A request by `T2` for an *X* lock cannot be granted immediately.

If a transaction `T1` holds an exclusive (*X*) lock on row `r`, a request from some distinct transaction `T2` for a lock of either type on `r` cannot be granted immediately. Instead, transaction `T2` has to wait for transaction `T1` to release its lock on row `r`.

Intention Locks

`InnoDB` supports *multiple granularity locking* which permits coexistence of row-level locks and locks on entire tables. To make locking at multiple granularity levels practical, additional types of locks called intention locks are used. Intention locks are table-level locks in `InnoDB` that indicate which type of lock (shared or exclusive) a transaction requires later for a row in that table. There are two types of intention locks used in `InnoDB` (assume that transaction `T` has requested a lock of the indicated type on table `t`):

- Intention shared (*IS*): Transaction T intends to set *S* locks on individual rows in table t.
- Intention exclusive (*IX*): Transaction T intends to set *X* locks on those rows.

For example, SELECT ... FOR SHARE sets an *IS* lock and SELECT ... FOR UPDATE sets an *IX* lock.

The intention locking protocol is as follows:

- Before a transaction can acquire an *S* lock on a row in table t, it must first acquire an *IS* or stronger lock on t.
- Before a transaction can acquire an *X* lock on a row, it must first acquire an *IX* lock on t.

These rules can be conveniently summarized by means of the following *lock type compatibility matrix*.

	X	***IX***	***S***	***IS***
X	Conflict	Conflict	Conflict	Conflict
IX	Conflict	Compatible	Conflict	Compatible
S	Conflict	Conflict	Compatible	Compatible
IS	Conflict	Compatible	Compatible	Compatible

A lock is granted to a requesting transaction if it is compatible with existing locks, but not if it conflicts with existing locks. A transaction waits until the conflicting existing lock is released. If a lock request conflicts with an existing lock and cannot be granted because it would cause deadlock, an error occurs.

Thus, intention locks do not block anything except full table requests (for example, LOCK TABLES ... WRITE). The main purpose of *IX* and *IS* locks is to show that someone is locking a row, or going to lock a row in the table.

Transaction data for an intention lock appears similar to the following in SHOW ENGINE INNODB STATUS and InnoDB monitor output:

```
TABLE LOCK table `test`.`t` trx id 10080 lock mode IX
```

Record Locks

A record lock is a lock on an index record. For example, SELECT c1 FROM t WHERE c1 = 10 FOR UPDATE; prevents any other transaction from inserting, updating, or deleting rows where the value of t.c1 is 10.

Record locks always lock index records, even if a table is defined with no indexes. For such cases, InnoDB creates a hidden clustered index and uses this index for record locking. See Section 15.8.2.1, “Clustered and Secondary Indexes”.

Transaction data for a record lock appears similar to the following in SHOW ENGINE INNODB STATUS and InnoDB monitor output:

```
RECORD LOCKS space id 58 page no 3 n bits 72 index `PRIMARY` of table `test`.`t`
trx id 10078 lock_mode X locks rec but not gap
Record lock, heap no 2 PHYSICAL RECORD: n_fields 3; compact format; info bits 0
 0: len 4; hex 8000000a; asc     ;;
 1: len 6; hex 00000000274f; asc     'O;;
 2: len 7; hex b60000019d0110; asc        ;;
```

Gap Locks

A gap lock is a lock on a gap between index records, or a lock on the gap before the first or after the last index record. For example, SELECT c1 FROM t WHERE c1 BETWEEN 10 and 20 FOR

`UPDATE;` prevents other transactions from inserting a value of `15` into column `t.c1`, whether or not there was already any such value in the column, because the gaps between all existing values in the range are locked.

A gap might span a single index value, multiple index values, or even be empty.

Gap locks are part of the tradeoff between performance and concurrency, and are used in some transaction isolation levels and not others.

Gap locking is not needed for statements that lock rows using a unique index to search for a unique row. (This does not include the case that the search condition includes only some columns of a multiple-column unique index; in that case, gap locking does occur.) For example, if the `id` column has a unique index, the following statement uses only an index-record lock for the row having `id` value 100 and it does not matter whether other sessions insert rows in the preceding gap:

```
SELECT * FROM child WHERE id = 100;
```

If `id` is not indexed or has a nonunique index, the statement does lock the preceding gap.

It is also worth noting here that conflicting locks can be held on a gap by different transactions. For example, transaction A can hold a shared gap lock (gap S-lock) on a gap while transaction B holds an exclusive gap lock (gap X-lock) on the same gap. The reason conflicting gap locks are allowed is that if a record is purged from an index, the gap locks held on the record by different transactions must be merged.

Gap locks in `InnoDB` are “purely inhibitive”, which means they only stop other transactions from inserting to the gap. They do not prevent different transactions from taking gap locks on the same gap. Thus, a gap X-lock has the same effect as a gap S-lock.

Gap locking can be disabled explicitly. This occurs if you change the transaction isolation level to `READ COMMITTED`. Under these circumstances, gap locking is disabled for searches and index scans and is used only for foreign-key constraint checking and duplicate-key checking.

There are also other effects of using the `READ COMMITTED` isolation level. Record locks for nonmatching rows are released after MySQL has evaluated the `WHERE` condition. For `UPDATE` statements, `InnoDB` does a “semi-consistent” read, such that it returns the latest committed version to MySQL so that MySQL can determine whether the row matches the `WHERE` condition of the `UPDATE`.

Next-Key Locks

A next-key lock is a combination of a record lock on the index record and a gap lock on the gap before the index record.

`InnoDB` performs row-level locking in such a way that when it searches or scans a table index, it sets shared or exclusive locks on the index records it encounters. Thus, the row-level locks are actually index-record locks. A next-key lock on an index record also affects the “gap” before that index record. That is, a next-key lock is an index-record lock plus a gap lock on the gap preceding the index record. If one session has a shared or exclusive lock on record `R` in an index, another session cannot insert a new index record in the gap immediately before `R` in the index order.

Suppose that an index contains the values 10, 11, 13, and 20. The possible next-key locks for this index cover the following intervals, where a round bracket denotes exclusion of the interval endpoint and a square bracket denotes inclusion of the endpoint:

```
(negative infinity, 10]
(10, 11]
(11, 13]
(13, 20]
(20, positive infinity)
```

For the last interval, the next-key lock locks the gap above the largest value in the index and the “supremum” pseudo-record having a value higher than any value actually in the index. The supremum is not a real index record, so, in effect, this next-key lock locks only the gap following the largest index value.

By default, `InnoDB` operates in `REPEATABLE READ` transaction isolation level. In this case, `InnoDB` uses next-key locks for searches and index scans, which prevents phantom rows (see Section 15.5.4, “Phantom Rows”).

Transaction data for a next-key lock appears similar to the following in `SHOW ENGINE INNODB STATUS` and InnoDB monitor output:

```
RECORD LOCKS space id 58 page no 3 n bits 72 index `PRIMARY` of table `test`.`t`
trx id 10080 lock_mode X
Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0
 0: len 8; hex 73757072656d756d; asc supremum;;

Record lock, heap no 2 PHYSICAL RECORD: n_fields 3; compact format; info bits 0
 0: len 4; hex 8000000a; asc     ;;
 1: len 6; hex 00000000274f; asc     'O;;
 2: len 7; hex b60000019d0110; asc        ;;
```

Insert Intention Locks

An insert intention lock is a type of gap lock set by `INSERT` operations prior to row insertion. This lock signals the intent to insert in such a way that multiple transactions inserting into the same index gap need not wait for each other if they are not inserting at the same position within the gap. Suppose that there are index records with values of 4 and 7. Separate transactions that attempt to insert values of 5 and 6, respectively, each lock the gap between 4 and 7 with insert intention locks prior to obtaining the exclusive lock on the inserted row, but do not block each other because the rows are nonconflicting.

The following example demonstrates a transaction taking an insert intention lock prior to obtaining an exclusive lock on the inserted record. The example involves two clients, A and B.

Client A creates a table containing two index records (90 and 102) and then starts a transaction that places an exclusive lock on index records with an ID greater than 100. The exclusive lock includes a gap lock before record 102:

```
mysql> CREATE TABLE child (id int(11) NOT NULL, PRIMARY KEY(id)) ENGINE=InnoDB;
mysql> INSERT INTO child (id) values (90),(102);

mysql> START TRANSACTION;
mysql> SELECT * FROM child WHERE id > 100 FOR UPDATE;
+-----+
| id  |
+-----+
| 102 |
+-----+
```

Client B begins a transaction to insert a record into the gap. The transaction takes an insert intention lock while it waits to obtain an exclusive lock.

```
mysql> START TRANSACTION;
mysql> INSERT INTO child (id) VALUES (101);
```

Transaction data for an insert intention lock appears similar to the following in `SHOW ENGINE INNODB STATUS` and InnoDB monitor output:

```
RECORD LOCKS space id 31 page no 3 n bits 72 index `PRIMARY` of table `test`.`child`
trx id 8731 lock_mode X locks gap before rec insert intention waiting
Record lock, heap no 3 PHYSICAL RECORD: n_fields 3; compact format; info bits 0
 0: len 4; hex 80000066; asc    f;;
 1: len 6; hex 000000002215; asc     " ;;
```

```
2: len 7; hex 9000000172011c; asc     r  ;;...
```

AUTO-INC Locks

An `AUTO-INC` lock is a special table-level lock taken by transactions inserting into tables with `AUTO_INCREMENT` columns. In the simplest case, if one transaction is inserting values into the table, any other transactions must wait to do their own inserts into that table, so that rows inserted by the first transaction receive consecutive primary key values.

The `innodb_autoinc_lock_mode` configuration option controls the algorithm used for auto-increment locking. It allows you to choose how to trade off between predictable sequences of auto-increment values and maximum concurrency for insert operations.

For more information, see Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”.

Predicate Locks for Spatial Indexes

`InnoDB` supports `SPATIAL` indexing of columns containing spatial columns (see Section 11.5.8, “Optimizing Spatial Analysis”).

To handle locking for operations involving `SPATIAL` indexes, next-key locking does not work well to support `REPEATABLE READ` or `SERIALIZABLE` transaction isolation levels. There is no absolute ordering concept in multidimensional data, so it is not clear which is the “next” key.

To enable support of isolation levels for tables with `SPATIAL` indexes, `InnoDB` uses predicate locks. A `SPATIAL` index contains minimum bounding rectangle (MBR) values, so `InnoDB` enforces consistent read on the index by setting a predicate lock on the MBR value used for a query. Other transactions cannot insert or modify a row that would match the query condition.

15.5.2 InnoDB Transaction Model

In the `InnoDB` transaction model, the goal is to combine the best properties of a multi-versioning database with traditional two-phase locking. `InnoDB` performs locking at the row level and runs queries as nonlocking consistent reads by default, in the style of Oracle. The lock information in `InnoDB` is stored space-efficiently so that lock escalation is not needed. Typically, several users are permitted to lock every row in `InnoDB` tables, or any random subset of the rows, without causing `InnoDB` memory exhaustion.

15.5.2.1 Transaction Isolation Levels

Transaction isolation is one of the foundations of database processing. Isolation is the I in the acronym ACID; the isolation level is the setting that fine-tunes the balance between performance and reliability, consistency, and reproducibility of results when multiple transactions are making changes and performing queries at the same time.

`InnoDB` offers all four transaction isolation levels described by the SQL:1992 standard: `READ UNCOMMITTED`, `READ COMMITTED`, `REPEATABLE READ`, and `SERIALIZABLE`. The default isolation level for `InnoDB` is `REPEATABLE READ`.

A user can change the isolation level for a single session or for all subsequent connections with the `SET TRANSACTION` statement. To set the server's default isolation level for all connections, use the `--transaction-isolation` option on the command line or in an option file. For detailed information about isolation levels and level-setting syntax, see Section 13.3.6, “SET TRANSACTION Syntax”.

`InnoDB` supports each of the transaction isolation levels described here using different locking strategies. You can enforce a high degree of consistency with the default `REPEATABLE READ` level, for operations on crucial data where ACID compliance is important. Or you can relax the consistency rules with `READ COMMITTED` or even `READ UNCOMMITTED`, in situations such as bulk reporting where precise consistency and repeatable results are less important than minimizing the amount of overhead for locking. `SERIALIZABLE` enforces even stricter rules than `REPEATABLE READ`, and is used mainly

in specialized situations, such as with XA transactions and for troubleshooting issues with concurrency and deadlocks.

The following list describes how MySQL supports the different transaction levels. The list goes from the most commonly used level to the least used.

- `REPEATABLE READ`

  This is the default isolation level for `InnoDB`. Consistent reads within the same transaction read the snapshot established by the first read. This means that if you issue several plain (nonlocking) `SELECT` statements within the same transaction, these `SELECT` statements are consistent also with respect to each other. See Section 15.5.2.3, “Consistent Nonlocking Reads”.

  For locking reads (`SELECT` with `FOR UPDATE` or `FOR SHARE`), `UPDATE`, and `DELETE` statements, locking depends on whether the statement uses a unique index with a unique search condition, or a range-type search condition.

  - For a unique index with a unique search condition, `InnoDB` locks only the index record found, not the gap before it.
  - For other search conditions, `InnoDB` locks the index range scanned, using gap locks or next-key locks to block insertions by other sessions into the gaps covered by the range. For information about gap locks and next-key locks, see Section 15.5.1, “InnoDB Locking”.

- `READ COMMITTED`

  Each consistent read, even within the same transaction, sets and reads its own fresh snapshot. For information about consistent reads, see Section 15.5.2.3, “Consistent Nonlocking Reads”.

  For locking reads (`SELECT` with `FOR UPDATE` or `FOR SHARE`), `UPDATE` statements, and `DELETE` statements, `InnoDB` locks only index records, not the gaps before them, and thus permits the free insertion of new records next to locked records. Gap locking is only used for foreign-key constraint checking and duplicate-key checking.

  Because gap locking is disabled, phantom problems may occur, as other sessions can insert new rows into the gaps. For information about phantoms, see Section 15.5.4, “Phantom Rows”.

  If you use `READ COMMITTED`, you *must* use row-based binary logging.

  Using `READ COMMITTED` has additional effects:

  - For `UPDATE` or `DELETE` statements, `InnoDB` holds locks only for rows that it updates or deletes. Record locks for nonmatching rows are released after MySQL has evaluated the `WHERE` condition. This greatly reduces the probability of deadlocks, but they can still happen.
  - For `UPDATE` statements, if a row is already locked, `InnoDB` performs a “semi-consistent” read, returning the latest committed version to MySQL so that MySQL can determine whether the row matches the `WHERE` condition of the `UPDATE`. If the row matches (must be updated), MySQL reads the row again and this time `InnoDB` either locks it or waits for a lock on it.

  Consider the following example, beginning with this table:

```
CREATE TABLE t (a INT NOT NULL, b INT) ENGINE = InnoDB;
INSERT INTO t VALUES (1,2),(2,3),(3,2),(4,3),(5,2);
COMMIT;
```

  In this case, table has no indexes, so searches and index scans use the hidden clustered index for record locking (see Section 15.8.2.1, “Clustered and Secondary Indexes”).

  Suppose that one client performs an `UPDATE` using these statements:

```
SET autocommit = 0;
UPDATE t SET b = 5 WHERE b = 3;
```

Suppose also that a second client performs an `UPDATE` by executing these statements following those of the first client:

```
SET autocommit = 0;
UPDATE t SET b = 4 WHERE b = 2;
```

As `InnoDB` executes each `UPDATE`, it first acquires an exclusive lock for each row, and then determines whether to modify it. If `InnoDB` does not modify the row, it releases the lock. Otherwise, `InnoDB` retains the lock until the end of the transaction. This affects transaction processing as follows.

When using the default `REPEATABLE READ` isolation level, the first `UPDATE` acquires x-locks and does not release any of them:

```
x-lock(1,2); retain x-lock
x-lock(2,3); update(2,3) to (2,5); retain x-lock
x-lock(3,2); retain x-lock
x-lock(4,3); update(4,3) to (4,5); retain x-lock
x-lock(5,2); retain x-lock
```

The second `UPDATE` blocks as soon as it tries to acquire any locks (because first update has retained locks on all rows), and does not proceed until the first `UPDATE` commits or rolls back:

```
x-lock(1,2); block and wait for first UPDATE to commit or roll back
```

If `READ COMMITTED` is used instead, the first `UPDATE` acquires x-locks and releases those for rows that it does not modify:

```
x-lock(1,2); unlock(1,2)
x-lock(2,3); update(2,3) to (2,5); retain x-lock
x-lock(3,2); unlock(3,2)
x-lock(4,3); update(4,3) to (4,5); retain x-lock
x-lock(5,2); unlock(5,2)
```

For the second `UPDATE`, `InnoDB` does a “semi-consistent” read, returning the latest committed version of each row to MySQL so that MySQL can determine whether the row matches the `WHERE` condition of the `UPDATE`:

```
x-lock(1,2); update(1,2) to (1,4); retain x-lock
x-lock(2,3); unlock(2,3)
x-lock(3,2); update(3,2) to (3,4); retain x-lock
x-lock(4,3); unlock(4,3)
x-lock(5,2); update(5,2) to (5,4); retain x-lock
```

- `READ UNCOMMITTED`

  `SELECT` statements are performed in a nonlocking fashion, but a possible earlier version of a row might be used. Thus, using this isolation level, such reads are not consistent. This is also called a dirty read. Otherwise, this isolation level works like `READ COMMITTED`.

- `SERIALIZABLE`

  This level is like `REPEATABLE READ`, but `InnoDB` implicitly converts all plain `SELECT` statements to `SELECT ... FOR SHARE` if `autocommit` is disabled. If `autocommit` is enabled, the `SELECT` is its own transaction. It therefore is known to be read only and can be serialized if performed as a

consistent (nonlocking) read and need not block for other transactions. (To force a plain `SELECT` to block if other transactions have modified the selected rows, disable `autocommit`.)

15.5.2.2 autocommit, Commit, and Rollback

In `InnoDB`, all user activity occurs inside a transaction. If `autocommit` mode is enabled, each SQL statement forms a single transaction on its own. By default, MySQL starts the session for each new connection with `autocommit` enabled, so MySQL does a commit after each SQL statement if that statement did not return an error. If a statement returns an error, the commit or rollback behavior depends on the error. See Section 15.20.4, “InnoDB Error Handling”.

A session that has `autocommit` enabled can perform a multiple-statement transaction by starting it with an explicit `START TRANSACTION` or `BEGIN` statement and ending it with a `COMMIT` or `ROLLBACK` statement. See Section 13.3.1, “START TRANSACTION, COMMIT, and ROLLBACK Syntax”.

If `autocommit` mode is disabled within a session with `SET autocommit = 0`, the session always has a transaction open. A `COMMIT` or `ROLLBACK` statement ends the current transaction and a new one starts.

If a session that has `autocommit` disabled ends without explicitly committing the final transaction, MySQL rolls back that transaction.

Some statements implicitly end a transaction, as if you had done a `COMMIT` before executing the statement. For details, see Section 13.3.3, “Statements That Cause an Implicit Commit”.

A `COMMIT` means that the changes made in the current transaction are made permanent and become visible to other sessions. A `ROLLBACK` statement, on the other hand, cancels all modifications made by the current transaction. Both `COMMIT` and `ROLLBACK` release all `InnoDB` locks that were set during the current transaction.

Grouping DML Operations with Transactions

By default, connection to the MySQL server begins with autocommit mode enabled, which automatically commits every SQL statement as you execute it. This mode of operation might be unfamiliar if you have experience with other database systems, where it is standard practice to issue a sequence of DML statements and commit them or roll them back all together.

To use multiple-statement transactions, switch autocommit off with the SQL statement `SET autocommit = 0` and end each transaction with `COMMIT` or `ROLLBACK` as appropriate. To leave autocommit on, begin each transaction with `START TRANSACTION` and end it with `COMMIT` or `ROLLBACK`. The following example shows two transactions. The first is committed; the second is rolled back.

```
shell> mysql test

mysql> CREATE TABLE customer (a INT, b CHAR (20), INDEX (a));
Query OK, 0 rows affected (0.00 sec)
mysql> -- Do a transaction with autocommit turned on.
mysql> START TRANSACTION;
Query OK, 0 rows affected (0.00 sec)
mysql> INSERT INTO customer VALUES (10, 'Heikki');
Query OK, 1 row affected (0.00 sec)
mysql> COMMIT;
Query OK, 0 rows affected (0.00 sec)
mysql> -- Do another transaction with autocommit turned off.
mysql> SET autocommit=0;
Query OK, 0 rows affected (0.00 sec)
mysql> INSERT INTO customer VALUES (15, 'John');
Query OK, 1 row affected (0.00 sec)
mysql> INSERT INTO customer VALUES (20, 'Paul');
Query OK, 1 row affected (0.00 sec)
mysql> DELETE FROM customer WHERE b = 'Heikki';
Query OK, 1 row affected (0.00 sec)
```

```
mysql> -- Now we undo those last 2 inserts and the delete.
mysql> ROLLBACK;
Query OK, 0 rows affected (0.00 sec)
mysql> SELECT * FROM customer;
+------+--------+
| a    | b      |
+------+--------+
|   10 | Heikki |
+------+--------+
1 row in set (0.00 sec)
mysql>
```

Transactions in Client-Side Languages

In APIs such as PHP, Perl DBI, JDBC, ODBC, or the standard C call interface of MySQL, you can send transaction control statements such as `COMMIT` to the MySQL server as strings just like any other SQL statements such as `SELECT` or `INSERT`. Some APIs also offer separate special transaction commit and rollback functions or methods.

15.5.2.3 Consistent Nonlocking Reads

A consistent read means that `InnoDB` uses multi-versioning to present to a query a snapshot of the database at a point in time. The query sees the changes made by transactions that committed before that point of time, and no changes made by later or uncommitted transactions. The exception to this rule is that the query sees the changes made by earlier statements within the same transaction. This exception causes the following anomaly: If you update some rows in a table, a `SELECT` sees the latest version of the updated rows, but it might also see older versions of any rows. If other sessions simultaneously update the same table, the anomaly means that you might see the table in a state that never existed in the database.

If the transaction isolation level is `REPEATABLE READ` (the default level), all consistent reads within the same transaction read the snapshot established by the first such read in that transaction. You can get a fresher snapshot for your queries by committing the current transaction and after that issuing new queries.

With `READ COMMITTED` isolation level, each consistent read within a transaction sets and reads its own fresh snapshot.

Consistent read is the default mode in which `InnoDB` processes `SELECT` statements in `READ COMMITTED` and `REPEATABLE READ` isolation levels. A consistent read does not set any locks on the tables it accesses, and therefore other sessions are free to modify those tables at the same time a consistent read is being performed on the table.

Suppose that you are running in the default `REPEATABLE READ` isolation level. When you issue a consistent read (that is, an ordinary `SELECT` statement), `InnoDB` gives your transaction a timepoint according to which your query sees the database. If another transaction deletes a row and commits after your timepoint was assigned, you do not see the row as having been deleted. Inserts and updates are treated similarly.

Note

The snapshot of the database state applies to `SELECT` statements within a transaction, not necessarily to DML statements. If you insert or modify some rows and then commit that transaction, a `DELETE` or `UPDATE` statement issued from another concurrent `REPEATABLE READ` transaction could affect those just-committed rows, even though the session could not query them. If a transaction does update or delete rows committed by a different transaction, those changes do become visible to the current transaction. For example, you might encounter a situation like the following:

```
SELECT COUNT(c1) FROM t1 WHERE c1 = 'xyz';
-- Returns 0: no rows match.
```

```
DELETE FROM t1 WHERE c1 = 'xyz';
-- Deletes several rows recently committed by other transaction.

SELECT COUNT(c2) FROM t1 WHERE c2 = 'abc';
-- Returns 0: no rows match.
UPDATE t1 SET c2 = 'cba' WHERE c2 = 'abc';
-- Affects 10 rows: another txn just committed 10 rows with 'abc' values.
SELECT COUNT(c2) FROM t1 WHERE c2 = 'cba';
-- Returns 10: this txn can now see the rows it just updated.
```

You can advance your timepoint by committing your transaction and then doing another `SELECT` or `START TRANSACTION WITH CONSISTENT SNAPSHOT`.

This is called *multi-versioned concurrency control*.

In the following example, session A sees the row inserted by B only when B has committed the insert and A has committed as well, so that the timepoint is advanced past the commit of B.

```
             Session A              Session B

           SET autocommit=0;      SET autocommit=0;
time
|          SELECT * FROM t;
|          empty set
|                                 INSERT INTO t VALUES (1, 2);
|
v          SELECT * FROM t;
           empty set
                                  COMMIT;

           SELECT * FROM t;
           empty set

           COMMIT;

           SELECT * FROM t;
           ---------------------
           |    1    |    2    |
           ---------------------
```

If you want to see the “freshest” state of the database, use either the `READ COMMITTED` isolation level or a locking read:

```
SELECT * FROM t FOR SHARE;
```

With `READ COMMITTED` isolation level, each consistent read within a transaction sets and reads its own fresh snapshot. With `FOR SHARE`, a locking read occurs instead: A `SELECT` blocks until the transaction containing the freshest rows ends (see Section 15.5.2.4, “Locking Reads”).

Consistent read does not work over certain DDL statements:

- Consistent read does not work over `DROP TABLE`, because MySQL cannot use a table that has been dropped and `InnoDB` destroys the table.
- Consistent read does not work over `ALTER TABLE`, because that statement makes a temporary copy of the original table and deletes the original table when the temporary copy is built. When you reissue a consistent read within a transaction, rows in the new table are not visible because those rows did not exist when the transaction's snapshot was taken. In this case, the transaction returns an error: `ER_TABLE_DEF_CHANGED`, “Table definition has changed, please retry transaction”.

The type of read varies for selects in clauses like `INSERT INTO ... SELECT`, `UPDATE ... (SELECT)`, and `CREATE TABLE ... SELECT` that do not specify `FOR UPDATE` or `FOR SHARE`:

- By default, `InnoDB` uses stronger locks and the `SELECT` part acts like `READ COMMITTED`, where each consistent read, even within the same transaction, sets and reads its own fresh snapshot.

- To use a consistent read in such cases, set the isolation level of the transaction to `READ UNCOMMITTED`, `READ COMMITTED`, or `REPEATABLE READ` (that is, anything other than `SERIALIZABLE`). In this case, no locks are set on rows read from the selected table.

15.5.2.4 Locking Reads

If you query data and then insert or update related data within the same transaction, the regular `SELECT` statement does not give enough protection. Other transactions can update or delete the same rows you just queried. `InnoDB` supports two types of locking reads that offer extra safety:

- `SELECT ... FOR SHARE`

  Sets a shared mode lock on any rows that are read. Other sessions can read the rows, but cannot modify them until your transaction commits. If any of these rows were changed by another transaction that has not yet committed, your query waits until that transaction ends and then uses the latest values.

Note

`SELECT ... FOR SHARE` is a replacement for `SELECT ... LOCK IN SHARE MODE`, but `LOCK IN SHARE MODE` remains available for backward compatibility. The statements are equivalent. However, `FOR SHARE` supports `OF table_name`, `NOWAIT`, and `SKIP LOCKED` options. See Locking Read Concurrency with NOWAIT and SKIP LOCKED.

- `SELECT ... FOR UPDATE`

  For index records the search encounters, locks the rows and any associated index entries, the same as if you issued an `UPDATE` statement for those rows. Other transactions are blocked from updating those rows, from doing `SELECT ... FOR SHARE`, or from reading the data in certain transaction isolation levels. Consistent reads ignore any locks set on the records that exist in the read view. (Old versions of a record cannot be locked; they are reconstructed by applying undo logs on an in-memory copy of the record.)

These clauses are primarily useful when dealing with tree-structured or graph-structured data, either in a single table or split across multiple tables. You traverse edges or tree branches from one place to another, while reserving the right to come back and change any of these “pointer” values.

All locks set by `FOR SHARE` and `FOR UPDATE` queries are released when the transaction is committed or rolled back.

Note

Locking of rows for update using `SELECT FOR UPDATE` only applies when autocommit is disabled (either by beginning transaction with `START TRANSACTION` or by setting `autocommit` to 0. If autocommit is enabled, the rows matching the specification are not locked.

Locking Read Examples

Suppose that you want to insert a new row into a table `child`, and make sure that the child row has a parent row in table `parent`. Your application code can ensure referential integrity throughout this sequence of operations.

First, use a consistent read to query the table `PARENT` and verify that the parent row exists. Can you safely insert the child row to table `CHILD`? No, because some other session could delete the parent row in the moment between your `SELECT` and your `INSERT`, without you being aware of it.

To avoid this potential issue, perform the `SELECT` using `FOR SHARE`:

```
SELECT * FROM parent WHERE NAME = 'Jones' FOR SHARE;
```

After the `FOR SHARE` query returns the parent `'Jones'`, you can safely add the child record to the `CHILD` table and commit the transaction. Any transaction that tries to acquire an exclusive lock in the applicable row in the `PARENT` table waits until you are finished, that is, until the data in all tables is in a consistent state.

For another example, consider an integer counter field in a table `CHILD_CODES`, used to assign a unique identifier to each child added to table `CHILD`. Do not use either consistent read or a shared mode read to read the present value of the counter, because two users of the database could see the same value for the counter, and a duplicate-key error occurs if two transactions attempt to add rows with the same identifier to the `CHILD` table.

Here, `FOR SHARE` is not a good solution because if two users read the counter at the same time, at least one of them ends up in deadlock when it attempts to update the counter.

To implement reading and incrementing the counter, first perform a locking read of the counter using `FOR UPDATE`, and then increment the counter. For example:

```
SELECT counter_field FROM child_codes FOR UPDATE;
UPDATE child_codes SET counter_field = counter_field + 1;
```

A `SELECT ... FOR UPDATE` reads the latest available data, setting exclusive locks on each row it reads. Thus, it sets the same locks a searched SQL `UPDATE` would set on the rows.

The preceding description is merely an example of how `SELECT ... FOR UPDATE` works. In MySQL, the specific task of generating a unique identifier actually can be accomplished using only a single access to the table:

```
UPDATE child_codes SET counter_field = LAST_INSERT_ID(counter_field + 1);
SELECT LAST_INSERT_ID();
```

The `SELECT` statement merely retrieves the identifier information (specific to the current connection). It does not access any table.

Locking Read Concurrency with NOWAIT and SKIP LOCKED

If a row is locked by a transaction, a `SELECT ... FOR UPDATE` or `SELECT ... FOR SHARE` transaction that requests the same locked row must wait until the blocking transaction releases the row lock. This behavior prevents transactions from updating or deleting rows that are queried for updates by other transactions. However, waiting for a row lock to be released is not necessary if you want the query to return immediately when a requested row is locked, or if excluding locked rows from the result set is acceptable.

To avoid waiting for other transactions to release row locks, `NO WAIT` and `SKIP LOCKED` options may be used with `SELECT ... FOR UPDATE` or `SELECT ... FOR SHARE` locking read statements.

- `NOWAIT`

  A locking read that uses `NOWAIT` never waits to acquire a row lock. The query executes immediately, failing with an error if a requested row is locked.

- `SKIP LOCKED`

  A locking read that uses `SKIP LOCKED` never waits to acquire a row lock. The query executes immediately, removing locked rows from the result set.

Note

Queries that skip locked rows return an inconsistent view of the data. `SKIP LOCKED` is therefore not suitable for general transactional work. However, it may be used to avoid lock contention when multiple sessions access the same queue-like table.

`NO WAIT` and `SKIP LOCKED` only apply to row-level locks.

Statements that use `NO WAIT` or `SKIP LOCKED` are unsafe for statement based replication.

The following example demonstrates `NOWAIT` and `SKIP LOCKED`. Session 1 starts a transaction that takes a row lock on a single record. Session 2 attempts a locking read on the same record using the `NOWAIT` option. Because the requested row is locked by Session 1, the locking read returns immediately with an error. In Session 3, the locking read with `SKIP LOCKED` returns the requested rows except for the row that is locked by Session 1.

```
# Session 1:

mysql> CREATE TABLE t (i INT, PRIMARY KEY (i)) ENGINE = InnoDB;

mysql> INSERT INTO t (i) VALUES(1),(2),(3);

mysql> START TRANSACTION;

mysql> SELECT * FROM t WHERE i = 2 FOR UPDATE;
+---+
| i |
+---+
| 2 |
+---+

# Session 2:

mysql> START TRANSACTION;

mysql> SELECT * FROM t WHERE i = 2 FOR UPDATE NOWAIT;
ERROR 3572 (HY000): Do not wait for lock.

# Session 3:

mysql> START TRANSACTION;

mysql> SELECT * FROM t FOR UPDATE SKIP LOCKED;
+---+
| i |
+---+
| 1 |
| 3 |
+---+
```

15.5.3 Locks Set by Different SQL Statements in InnoDB

A locking read, an `UPDATE`, or a `DELETE` generally set record locks on every index record that is scanned in the processing of the SQL statement. It does not matter whether there are `WHERE` conditions in the statement that would exclude the row. `InnoDB` does not remember the exact `WHERE` condition, but only knows which index ranges were scanned. The locks are normally next-key locks that also block inserts into the “gap” immediately before the record. However, gap locking can be disabled explicitly, which causes next-key locking not to be used. For more information, see Section 15.5.1, “InnoDB Locking”. The transaction isolation level also can affect which locks are set; see Section 15.5.2.1, “Transaction Isolation Levels”.

If a secondary index is used in a search and index record locks to be set are exclusive, `InnoDB` also retrieves the corresponding clustered index records and sets locks on them.

Differences between shared and exclusive locks are described in Section 15.5.1, “InnoDB Locking”.

If you have no indexes suitable for your statement and MySQL must scan the entire table to process the statement, every row of the table becomes locked, which in turn blocks all inserts by other users to the table. It is important to create good indexes so that your queries do not unnecessarily scan many rows.

For `SELECT ... FOR UPDATE` or `SELECT ... FOR SHARE`, locks are acquired for scanned rows, and expected to be released for rows that do not qualify for inclusion in the result set (for example, if they do not meet the criteria given in the `WHERE` clause). However, in some cases, rows might not be unlocked immediately because the relationship between a result row and its original source is lost during query execution. For example, in a `UNION`, scanned (and locked) rows from a table might be inserted into a temporary table before evaluation whether they qualify for the result set. In this circumstance, the relationship of the rows in the temporary table to the rows in the original table is lost and the latter rows are not unlocked until the end of query execution.

`InnoDB` sets specific types of locks as follows.

- `SELECT ... FROM` is a consistent read, reading a snapshot of the database and setting no locks unless the transaction isolation level is set to `SERIALIZABLE`. For `SERIALIZABLE` level, the search sets shared next-key locks on the index records it encounters. However, only an index record lock is required for statements that lock rows using a unique index to search for a unique row.

- `SELECT ... FROM ... FOR SHARE` sets shared next-key locks on all index records the search encounters. However, only an index record lock is required for statements that lock rows using a unique index to search for a unique row.

- `SELECT ... FROM ... FOR UPDATE` sets an exclusive next-key lock on every record the search encounters. However, only an index record lock is required for statements that lock rows using a unique index to search for a unique row.

  For index records the search encounters, `SELECT ... FROM ... FOR UPDATE` blocks other sessions from doing `SELECT ... FROM ... FOR SHARE` or from reading in certain transaction isolation levels. Consistent reads ignore any locks set on the records that exist in the read view.

- `UPDATE ... WHERE ...` sets an exclusive next-key lock on every record the search encounters. However, only an index record lock is required for statements that lock rows using a unique index to search for a unique row.

- When `UPDATE` modifies a clustered index record, implicit locks are taken on affected secondary index records. The `UPDATE` operation also takes shared locks on affected secondary index records when performing duplicate check scans prior to inserting new secondary index records, and when inserting new secondary index records.

- `DELETE FROM ... WHERE ...` sets an exclusive next-key lock on every record the search encounters. However, only an index record lock is required for statements that lock rows using a unique index to search for a unique row.

- `INSERT` sets an exclusive lock on the inserted row. This lock is an index-record lock, not a next-key lock (that is, there is no gap lock) and does not prevent other sessions from inserting into the gap before the inserted row.

  Prior to inserting the row, a type of gap lock called an insert intention gap lock is set. This lock signals the intent to insert in such a way that multiple transactions inserting into the same index gap need not wait for each other if they are not inserting at the same position within the gap. Suppose that there are index records with values of 4 and 7. Separate transactions that attempt to insert values of 5 and 6 each lock the gap between 4 and 7 with insert intention locks prior to obtaining the exclusive lock on the inserted row, but do not block each other because the rows are nonconflicting.

  If a duplicate-key error occurs, a shared lock on the duplicate index record is set. This use of a shared lock can result in deadlock should there be multiple sessions trying to insert the same row if another session already has an exclusive lock. This can occur if another session deletes the row. Suppose that an `InnoDB` table `t1` has the following structure:

  ```
  CREATE TABLE t1 (i INT, PRIMARY KEY (i)) ENGINE = InnoDB;
  ```

  Now suppose that three sessions perform the following operations in order:

Session 1:

```
START TRANSACTION;
INSERT INTO t1 VALUES(1);
```

Session 2:

```
START TRANSACTION;
INSERT INTO t1 VALUES(1);
```

Session 3:

```
START TRANSACTION;
INSERT INTO t1 VALUES(1);
```

Session 1:

```
ROLLBACK;
```

The first operation by session 1 acquires an exclusive lock for the row. The operations by sessions 2 and 3 both result in a duplicate-key error and they both request a shared lock for the row. When session 1 rolls back, it releases its exclusive lock on the row and the queued shared lock requests for sessions 2 and 3 are granted. At this point, sessions 2 and 3 deadlock: Neither can acquire an exclusive lock for the row because of the shared lock held by the other.

A similar situation occurs if the table already contains a row with key value 1 and three sessions perform the following operations in order:

Session 1:

```
START TRANSACTION;
DELETE FROM t1 WHERE i = 1;
```

Session 2:

```
START TRANSACTION;
INSERT INTO t1 VALUES(1);
```

Session 3:

```
START TRANSACTION;
INSERT INTO t1 VALUES(1);
```

Session 1:

```
COMMIT;
```

The first operation by session 1 acquires an exclusive lock for the row. The operations by sessions 2 and 3 both result in a duplicate-key error and they both request a shared lock for the row. When session 1 commits, it releases its exclusive lock on the row and the queued shared lock requests for sessions 2 and 3 are granted. At this point, sessions 2 and 3 deadlock: Neither can acquire an exclusive lock for the row because of the shared lock held by the other.

- `INSERT ... ON DUPLICATE KEY UPDATE` differs from a simple `INSERT` in that an exclusive lock rather than a shared lock is placed on the row to be updated when a duplicate-key error occurs. An exclusive index-record lock is taken for a duplicate primary key value. An exclusive next-key lock is taken for a duplicate unique key value.

- `REPLACE` is done like an `INSERT` if there is no collision on a unique key. Otherwise, an exclusive next-key lock is placed on the row to be replaced.

- `INSERT INTO T SELECT ... FROM S WHERE ...` sets an exclusive index record lock (without a gap lock) on each row inserted into `T`. If the transaction isolation level is `READ COMMITTED`, `InnoDB` does the search on `S` as a consistent read (no locks). Otherwise, `InnoDB` sets shared next-key locks on rows from `S`. `InnoDB` has to set locks in the latter case: In roll-forward recovery from a backup, every SQL statement must be executed in exactly the same way it was done originally.

  `CREATE TABLE ... SELECT ...` performs the `SELECT` with shared next-key locks or as a consistent read, as for `INSERT ... SELECT`.

  When a `SELECT` is used in the constructs `REPLACE INTO t SELECT ... FROM s WHERE ...` or `UPDATE t ... WHERE col IN (SELECT ... FROM s ...)`, `InnoDB` sets shared next-key locks on rows from table `s`.

- While initializing a previously specified `AUTO_INCREMENT` column on a table, `InnoDB` sets an exclusive lock on the end of the index associated with the `AUTO_INCREMENT` column. In accessing the auto-increment counter, `InnoDB` uses a specific `AUTO-INC` table lock mode where the lock lasts only to the end of the current SQL statement, not to the end of the entire transaction. Other sessions cannot insert into the table while the `AUTO-INC` table lock is held; see Section 15.5.2, “InnoDB Transaction Model”.

  `InnoDB` fetches the value of a previously initialized `AUTO_INCREMENT` column without setting any locks.

- If a `FOREIGN KEY` constraint is defined on a table, any insert, update, or delete that requires the constraint condition to be checked sets shared record-level locks on the records that it looks at to check the constraint. `InnoDB` also sets these locks in the case where the constraint fails.

- `LOCK TABLES` sets table locks, but it is the higher MySQL layer above the `InnoDB` layer that sets these locks. `InnoDB` is aware of table locks if `innodb_table_locks = 1` (the default) and `autocommit = 0`, and the MySQL layer above `InnoDB` knows about row-level locks.

  Otherwise, `InnoDB`'s automatic deadlock detection cannot detect deadlocks where such table locks are involved. Also, because in this case the higher MySQL layer does not know about row-level locks, it is possible to get a table lock on a table where another session currently has row-level locks. However, this does not endanger transaction integrity, as discussed in Section 15.5.5.2, “Deadlock Detection and Rollback”. See also Section 15.8.1.7, “Limits on InnoDB Tables”.

15.5.4 Phantom Rows

The so-called *phantom* problem occurs within a transaction when the same query produces different sets of rows at different times. For example, if a `SELECT` is executed twice, but returns a row the second time that was not returned the first time, the row is a “phantom” row.

Suppose that there is an index on the `id` column of the `child` table and that you want to read and lock all rows from the table having an identifier value larger than 100, with the intention of updating some column in the selected rows later:

```
SELECT * FROM child WHERE id > 100 FOR UPDATE;
```

The query scans the index starting from the first record where `id` is bigger than 100. Let the table contain rows having `id` values of 90 and 102. If the locks set on the index records in the scanned range do not lock out inserts made in the gaps (in this case, the gap between 90 and 102), another session can insert a new row into the table with an `id` of 101. If you were to execute the same `SELECT` within the same transaction, you would see a new row with an `id` of 101 (a “phantom”) in the result set returned by the query. If we regard a set of rows as a data item, the new phantom child would violate the isolation principle of transactions that a transaction should be able to run so that the data it has read does not change during the transaction.

To prevent phantoms, `InnoDB` uses an algorithm called *next-key locking* that combines index-row locking with gap locking. `InnoDB` performs row-level locking in such a way that when it searches or scans a table index, it sets shared or exclusive locks on the index records it encounters. Thus, the row-level locks are actually index-record locks. In addition, a next-key lock on an index record also affects the "gap" before that index record. That is, a next-key lock is an index-record lock plus a gap lock on the gap preceding the index record. If one session has a shared or exclusive lock on record `R` in an index, another session cannot insert a new index record in the gap immediately before `R` in the index order.

When `InnoDB` scans an index, it can also lock the gap after the last record in the index. Just that happens in the preceding example: To prevent any insert into the table where `id` would be bigger than 100, the locks set by `InnoDB` include a lock on the gap following `id` value 102.

You can use next-key locking to implement a uniqueness check in your application: If you read your data in share mode and do not see a duplicate for a row you are going to insert, then you can safely insert your row and know that the next-key lock set on the successor of your row during the read prevents anyone meanwhile inserting a duplicate for your row. Thus, the next-key locking enables you to "lock" the nonexistence of something in your table.

Gap locking can be disabled as discussed in Section 15.5.1, "InnoDB Locking". This may cause phantom problems because other sessions can insert new rows into the gaps when gap locking is disabled.

15.5.5 Deadlocks in InnoDB

A deadlock is a situation where different transactions are unable to proceed because each holds a lock that the other needs. Because both transactions are waiting for a resource to become available, neither ever release the locks it holds.

A deadlock can occur when transactions lock rows in multiple tables (through statements such as `UPDATE` or `SELECT ... FOR UPDATE`), but in the opposite order. A deadlock can also occur when such statements lock ranges of index records and gaps, with each transaction acquiring some locks but not others due to a timing issue. For a deadlock example, see Section 15.5.5.1, "An InnoDB Deadlock Example".

To reduce the possibility of deadlocks, use transactions rather than `LOCK TABLES` statements; keep transactions that insert or update data small enough that they do not stay open for long periods of time; when different transactions update multiple tables or large ranges of rows, use the same order of operations (such as `SELECT ... FOR UPDATE`) in each transaction; create indexes on the columns used in `SELECT ... FOR UPDATE` and `UPDATE ... WHERE` statements. The possibility of deadlocks is not affected by the isolation level, because the isolation level changes the behavior of read operations, while deadlocks occur because of write operations. For more information about avoiding and recovering from deadlock conditions, see Section 15.5.5.3, "How to Minimize and Handle Deadlocks".

When deadlock detection is enabled (the default) and a deadlock does occur, `InnoDB` detects the condition and rolls back one of the transactions (the victim). If deadlock detection is disabled using the `innodb_deadlock_detect` configuration option, `InnoDB` relies on the `innodb_lock_wait_timeout` setting to roll back transactions in case of a deadlock. Thus, even if your application logic is correct, you must still handle the case where a transaction must be retried. To see the last deadlock in an `InnoDB` user transaction, use the `SHOW ENGINE INNODB STATUS` command. If frequent deadlocks highlight a problem with transaction structure or application error handling, run with the `innodb_print_all_deadlocks` setting enabled to print information about all deadlocks to the `mysqld` error log. For more information about how deadlocks are automatically detected and handled, see Section 15.5.5.2, "Deadlock Detection and Rollback".

15.5.5.1 An InnoDB Deadlock Example

The following example illustrates how an error can occur when a lock request would cause a deadlock. The example involves two clients, A and B.

First, client A creates a table containing one row, and then begins a transaction. Within the transaction, A obtains an *S* lock on the row by selecting it in share mode:

```
mysql> CREATE TABLE t (i INT) ENGINE = InnoDB;
Query OK, 0 rows affected (1.07 sec)

mysql> INSERT INTO t (i) VALUES(1);
Query OK, 1 row affected (0.09 sec)

mysql> START TRANSACTION;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT * FROM t WHERE i = 1 FOR SHARE;
+------+
| i    |
+------+
|    1 |
+------+
```

Next, client B begins a transaction and attempts to delete the row from the table:

```
mysql> START TRANSACTION;
Query OK, 0 rows affected (0.00 sec)

mysql> DELETE FROM t WHERE i = 1;
```

The delete operation requires an *X* lock. The lock cannot be granted because it is incompatible with the *S* lock that client A holds, so the request goes on the queue of lock requests for the row and client B blocks.

Finally, client A also attempts to delete the row from the table:

```
mysql> DELETE FROM t WHERE i = 1;
ERROR 1213 (40001): Deadlock found when trying to get lock;
try restarting transaction
```

Deadlock occurs here because client A needs an *X* lock to delete the row. However, that lock request cannot be granted because client B already has a request for an *X* lock and is waiting for client A to release its *S* lock. Nor can the *S* lock held by A be upgraded to an *X* lock because of the prior request by B for an *X* lock. As a result, `InnoDB` generates an error for one of the clients and releases its locks. The client returns this error:

```
ERROR 1213 (40001): Deadlock found when trying to get lock;
try restarting transaction
```

At that point, the lock request for the other client can be granted and it deletes the row from the table.

15.5.5.2 Deadlock Detection and Rollback

When deadlock detection is enabled (the default), `InnoDB` automatically detects transaction deadlocks and rolls back a transaction or transactions to break the deadlock. `InnoDB` tries to pick small transactions to roll back, where the size of a transaction is determined by the number of rows inserted, updated, or deleted.

`InnoDB` is aware of table locks if `innodb_table_locks = 1` (the default) and `autocommit = 0`, and the MySQL layer above it knows about row-level locks. Otherwise, `InnoDB` cannot detect deadlocks where a table lock set by a MySQL `LOCK TABLES` statement or a lock set by a storage engine other than `InnoDB` is involved. Resolve these situations by setting the value of the `innodb_lock_wait_timeout` system variable.

When `InnoDB` performs a complete rollback of a transaction, all locks set by the transaction are released. However, if just a single SQL statement is rolled back as a result of an error, some of the

locks set by the statement may be preserved. This happens because InnoDB stores row locks in a format such that it cannot know afterward which lock was set by which statement.

If a SELECT calls a stored function in a transaction, and a statement within the function fails, that statement rolls back. Furthermore, if ROLLBACK is executed after that, the entire transaction rolls back.

If the LATEST DETECTED DEADLOCK section of InnoDB Monitor output includes a message stating, "TOO DEEP OR LONG SEARCH IN THE LOCK TABLE WAITS-FOR GRAPH, WE WILL ROLL BACK FOLLOWING TRANSACTION," this indicates that the number of transactions on the wait-for list has reached a limit of 200. A wait-for list that exceeds 200 transactions is treated as a deadlock and the transaction attempting to check the wait-for list is rolled back. The same error may also occur if the locking thread must look at more than 1,000,000 locks owned by transactions on the wait-for list.

For techniques to organize database operations to avoid deadlocks, see Section 15.5.5, "Deadlocks in InnoDB".

Disabling Deadlock Detection

On high concurrency systems, deadlock detection can cause a slowdown when numerous threads wait for the same lock. At times, it may be more efficient to disable deadlock detection and rely on the innodb_lock_wait_timeout setting for transaction rollback when a deadlock occurs. Deadlock detection can be disabled using the innodb_deadlock_detect configuration option.

15.5.5.3 How to Minimize and Handle Deadlocks

This section builds on the conceptual information about deadlocks in Section 15.5.5.2, "Deadlock Detection and Rollback". It explains how to organize database operations to minimize deadlocks and the subsequent error handling required in applications.

Deadlocks are a classic problem in transactional databases, but they are not dangerous unless they are so frequent that you cannot run certain transactions at all. Normally, you must write your applications so that they are always prepared to re-issue a transaction if it gets rolled back because of a deadlock.

InnoDB uses automatic row-level locking. You can get deadlocks even in the case of transactions that just insert or delete a single row. That is because these operations are not really "atomic"; they automatically set locks on the (possibly several) index records of the row inserted or deleted.

You can cope with deadlocks and reduce the likelihood of their occurrence with the following techniques:

- At any time, issue the SHOW ENGINE INNODB STATUS command to determine the cause of the most recent deadlock. That can help you to tune your application to avoid deadlocks.
- If frequent deadlock warnings cause concern, collect more extensive debugging information by enabling the innodb_print_all_deadlocks configuration option. Information about each deadlock, not just the latest one, is recorded in the MySQL error log. Disable this option when you are finished debugging.
- Always be prepared to re-issue a transaction if it fails due to deadlock. Deadlocks are not dangerous. Just try again.
- Keep transactions small and short in duration to make them less prone to collision.
- Commit transactions immediately after making a set of related changes to make them less prone to collision. In particular, do not leave an interactive mysql session open for a long time with an uncommitted transaction.
- If you use locking reads (SELECT ... FOR UPDATE or SELECT ... FOR SHARE), try using a lower isolation level such as READ COMMITTED.

- When modifying multiple tables within a transaction, or different sets of rows in the same table, do those operations in a consistent order each time. Then transactions form well-defined queues and do not deadlock. For example, organize database operations into functions within your application, or call stored routines, rather than coding multiple similar sequences of `INSERT`, `UPDATE`, and `DELETE` statements in different places.

- Add well-chosen indexes to your tables. Then your queries need to scan fewer index records and consequently set fewer locks. Use `EXPLAIN SELECT` to determine which indexes the MySQL server regards as the most appropriate for your queries.

- Use less locking. If you can afford to permit a `SELECT` to return data from an old snapshot, do not add the clause `FOR UPDATE` or `FOR SHARE` to it. Using the `READ COMMITTED` isolation level is good here, because each consistent read within the same transaction reads from its own fresh snapshot.

- If nothing else helps, serialize your transactions with table-level locks. The correct way to use `LOCK TABLES` with transactional tables, such as `InnoDB` tables, is to begin a transaction with `SET autocommit = 0` (not `START TRANSACTION`) followed by `LOCK TABLES`, and to not call `UNLOCK TABLES` until you commit the transaction explicitly. For example, if you need to write to table `t1` and read from table `t2`, you can do this:

  ```
  SET autocommit=0;
  LOCK TABLES t1 WRITE, t2 READ, ...;
  ... do something with tables t1 and t2 here ...
  COMMIT;
  UNLOCK TABLES;
  ```

  Table-level locks prevent concurrent updates to the table, avoiding deadlocks at the expense of less responsiveness for a busy system.

- Another way to serialize transactions is to create an auxiliary “semaphore” table that contains just a single row. Have each transaction update that row before accessing other tables. In that way, all transactions happen in a serial fashion. Note that the `InnoDB` instant deadlock detection algorithm also works in this case, because the serializing lock is a row-level lock. With MySQL table-level locks, the timeout method must be used to resolve deadlocks.

15.6 InnoDB Configuration

This section provides configuration information and procedures for `InnoDB` initialization, startup, and various components and features of the `InnoDB` storage engine. For information about optimizing database operations for `InnoDB` tables, see Section 8.5, “Optimizing for InnoDB Tables”.

15.6.1 InnoDB Startup Configuration

The first decisions to make about `InnoDB` configuration involve the configuration of data files, log files, page size, and memory buffers. It is recommended that you define data file, log file, and page size configuration before creating the `InnoDB` instance. Modifying data file or log file configuration after the `InnoDB` instance is created may involve a non-trivial procedure, and page size can only be defined when the `InnoDB` instance is first initialized.

In addition to these topics, this section provides information about specifying `InnoDB` options in a configuration file, viewing `InnoDB` initialization information, and important storage considerations.

- Specifying Options in a MySQL Configuration File
- Viewing InnoDB Initialization Information
- Important Storage Considerations
- System Tablespace Data File Configuration

- InnoDB Log File Configuration
- InnoDB Undo Tablespace Configuration
- InnoDB Temporary Tablespace Configuration
- InnoDB Page Size Configuration
- InnoDB Memory Configuration

Specifying Options in a MySQL Configuration File

Because MySQL uses data file, log file, and page size configuration settings to initialize the `InnoDB` instance, it is recommended that you define these settings in a configuration file that MySQL reads at startup, prior to initializing `InnoDB` for the first time. `InnoDB` is initialized when the MySQL server is started, and the first initialization of `InnoDB` normally occurs the first time you start the MySQL server.

You can place `InnoDB` options in the `[mysqld]` group of any option file that your server reads when it starts. The locations of MySQL option files are described in Section 4.2.6, “Using Option Files”.

To make sure that `mysqld` reads options only from a specific file (and `mysqld-auto.cnf`), use the `--defaults-file` option as the first option on the command line when starting the server:

```
mysqld --defaults-file=path_to_configuration_file
```

Viewing InnoDB Initialization Information

To view `InnoDB` initialization information during startup, start `mysqld` from a command prompt. When `mysqld` is started from a command prompt, initialization information is printed to the console.

For example, on Windows, if `mysqld` is located in `C:\Program Files\MySQL\MySQL Server 8.0\bin`, start the MySQL server like this:

```
C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld" --console
```

On Unix-like systems, `mysqld` is located in the `bin` directory of your MySQL installation:

```
sell> bin/mysqld --user=mysql &
```

If you do not send server output to the console, check the error log after startup to see the initialization information `InnoDB` printed during the startup process.

For information about starting MySQL using other methods, see Section 2.9.5, “Starting and Stopping MySQL Automatically”.

Important Storage Considerations

Review the following storage-related considerations before proceeding with your startup configuration.

- In some cases, database performance improves if the data is not all placed on the same physical disk. Putting log files on a different disk from data is very often beneficial for performance. For example, you can place system tablespace data files and log files on different disks. You can also use raw disk partitions (raw devices) for `InnoDB` data files, which may speed up I/O. See Section 15.7.3, “Using Raw Disk Partitions for the System Tablespace”.

- `InnoDB` is a transaction-safe (ACID compliant) storage engine for MySQL that has commit, rollback, and crash-recovery capabilities to protect user data. **However, it cannot do so** if the underlying operating system or hardware does not work as advertised. Many operating systems or disk subsystems may delay or reorder write operations to improve performance. On some operating systems, the very `fsync()` system call that should wait until all unwritten data for a file has been

flushed might actually return before the data has been flushed to stable storage. Because of this, an operating system crash or a power outage may destroy recently committed data, or in the worst case, even corrupt the database because of write operations having been reordered. If data integrity is important to you, perform some “pull-the-plug” tests before using anything in production. On OS X 10.3 and higher, `InnoDB` uses a special `fcntl()` file flush method. Under Linux, it is advisable to **disable the write-back cache**.

On ATA/SATA disk drives, a command such `hdparm -W0 /dev/hda` may work to disable the write-back cache. **Beware that some drives or disk controllers may be unable to disable the write-back cache.**

- With regard to `InnoDB` recovery capabilities that protect user data, `InnoDB` uses a file flush technique involving a structure called the doublewrite buffer, which is enabled by default (`innodb_doublewrite=ON`). The doublewrite buffer adds safety to recovery following a crash or power outage, and improves performance on most varieties of Unix by reducing the need for `fsync()` operations. It is recommended that the `innodb_doublewrite` option remains enabled if you are concerned with data integrity or possible failures. For additional information about the doublewrite buffer, see Section 15.11.1, “InnoDB Disk I/O”.
- Before using NFS with `InnoDB`, review potential issues outlined in Using NFS with MySQL.

System Tablespace Data File Configuration

System tablespace data files are configured using the `innodb_data_file_path` and `innodb_data_home_dir` configuration options.

The `innodb_data_file_path` configuration option is used to configure the `InnoDB` system tablespace data files. The value of `innodb_data_file_path` should be a list of one or more data file specifications. If you name more than one data file, separate them by semicolon (`;`) characters:

```
innodb_data_file_path=datafile_spec1[;datafile_spec2]...
```

For example, the following setting explicitly creates a minimally sized system tablespace:

```
[mysqld]
innodb_data_file_path=ibdata1:12M:autoextend
```

This setting configures a single 12MB data file named `ibdata1` that is auto-extending. No location for the file is given, so by default, `InnoDB` creates it in the MySQL data directory.

Sizes are specified using `K`, `M`, or `G` suffix letters to indicate units of KB, MB, or GB.

A tablespace containing a fixed-size 50MB data file named `ibdata1` and a 50MB auto-extending file named `ibdata2` in the data directory can be configured like this:

```
[mysqld]
innodb_data_file_path=ibdata1:50M;ibdata2:50M:autoextend
```

The full syntax for a data file specification includes the file name, its size, and several optional attributes:

```
file_name:file_size[:autoextend[:max:max_file_size]]
```

The `autoextend` and `max` attributes can be used only for the last data file in the `innodb_data_file_path` line.

If you specify the `autoextend` option for the last data file, `InnoDB` extends the data file if it runs out of free space in the tablespace. The increment is 64MB at a time by default. To modify the increment, change the `innodb_autoextend_increment` system variable.

If the disk becomes full, you might want to add another data file on another disk. For tablespace reconfiguration instructions, see Section 15.7.1, “Resizing the InnoDB System Tablespace”.

InnoDB is not aware of the file system maximum file size, so be cautious on file systems where the maximum file size is a small value such as 2GB. To specify a maximum size for an auto-extending data file, use the max attribute following the autoextend attribute. Use the max attribute only in cases where constraining disk usage is of critical importance, because exceeding the maximum size causes a fatal error, possibly including a crash. The following configuration permits ibdata1 to grow up to a limit of 500MB:

```
[mysqld]
innodb_data_file_path=ibdata1:12M:autoextend:max:500M
```

InnoDB creates tablespace files in the MySQL data directory by default (datadir). To specify a location explicitly, use the innodb_data_home_dir option. For example, to create two files named ibdata1 and ibdata2 in a directory named myibdata, configure InnoDB like this:

```
[mysqld]
innodb_data_home_dir = /path/to/myibdata/
innodb_data_file_path=ibdata1:50M;ibdata2:50M:autoextend
```

Note

A trailing slash is required when specifying a value for innodb_data_home_dir.

InnoDB does not create directories, so make sure that the myibdata directory exists before you start the server. Use the Unix or DOS mkdir command to create any necessary directories.

Make sure that the MySQL server has the proper access rights to create files in the data directory. More generally, the server must have access rights in any directory where it needs to create data files.

InnoDB forms the directory path for each data file by textually concatenating the value of innodb_data_home_dir to the data file name. If the innodb_data_home_dir option is not specified in my.cnf at all, the default value is the “dot” directory ./, which means the MySQL data directory. (The MySQL server changes its current working directory to its data directory when it begins executing.)

If you specify innodb_data_home_dir as an empty string, you can specify absolute paths for the data files listed in the innodb_data_file_path value. The following example is equivalent to the preceding one:

```
[mysqld]
innodb_data_home_dir =
innodb_data_file_path=/path/to/myibdata/ibdata1:50M;/path/to/myibdata/ibdata2:50M:autoextend
```

InnoDB Log File Configuration

By default, InnoDB creates two 48MB log files in the MySQL data directory (datadir) named ib_logfile0 and ib_logfile1.

The following options can be used to modify the default configuration:

- innodb_log_group_home_dir defines directory path to the InnoDB log files (the redo logs). If this option is not configured, InnoDB log files are created in the MySQL data directory (datadir).

  You might use this option to place InnoDB log files in a different physical storage location than InnoDB data files to avoid potential I/O resource conflicts. For example:

```
[mysqld]
innodb_log_group_home_dir = /dr3/iblogs
```

Note

`InnoDB` does not create directories, so make sure that the log directory exists before you start the server. Use the Unix or DOS `mkdir` command to create any necessary directories.

Make sure that the MySQL server has the proper access rights to create files in the log directory. More generally, the server must have access rights in any directory where it needs to create log files.

- `innodb_log_files_in_group` defines the number of log files in the log group. The default and recommended value is 2.
- `innodb_log_file_size` defines the size in bytes of each log file in the log group. The combined size of log files (`innodb_log_file_size` * `innodb_log_files_in_group`) cannot exceed a maximum value that is slightly less than 512GB. A pair of 255 GB log files, for example, approaches the limit but does not exceed it. The default log file size is 48MB. Generally, the combined size of the log files should be large enough that the server can smooth out peaks and troughs in workload activity, which often means that there is enough redo log space to handle more than an hour of write activity. The larger the value, the less checkpoint flush activity is needed in the buffer pool, saving disk I/O. For additional information, see Section 8.5.4, “Optimizing InnoDB Redo Logging”.

InnoDB Undo Tablespace Configuration

By default, `InnoDB` undo logs are part of the system tablespace. However, you can choose to store `InnoDB` undo logs in one or more separate undo tablespaces, typically on a different storage device.

The `innodb_undo_directory` configuration option defines the path where `InnoDB` creates separate tablespaces for the undo logs. This option is typically used in conjunction with the `innodb_rollback_segments` and `innodb_undo_tablespaces` options, which determine the disk layout of the undo logs outside the system tablespace.

For more information, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

InnoDB Temporary Tablespace Configuration

By default, `InnoDB` creates a single auto-extending temporary tablespace data file named `ibtmp1` that is slightly larger than 12MB in the `innodb_data_home_dir` directory. The default temporary tablespace data file configuration can be modified at startup using the `innodb_temp_data_file_path` configuration option.

The `innodb_temp_data_file_path` option specifies the path, file name, and file size for `InnoDB` temporary tablespace data files. The full directory path for a file is formed by concatenating `innodb_data_home_dir` to the path specified by `innodb_temp_data_file_path`. File size is specified in KB, MB, or GB (1024MB) by appending K, M, or G to the size value. The sum of the sizes of the files must be slightly larger than 12MB.

The `innodb_data_home_dir` default value is the MySQL data directory (`datadir`).

InnoDB Page Size Configuration

The `innodb_page_size` option specifies the page size for all `InnoDB` tablespaces in a MySQL instance. This value is set when the instance is created and remains constant afterward. Valid values are 64k, 32k, 16k (the default), 8k, and 4k. Alternatively, you can specify page size in bytes (65536, 32768, 16384, 8192, 4096).

The default page size of 16k is appropriate for a wide range of workloads, particularly for queries involving table scans and DML operations involving bulk updates. Smaller page sizes might be more efficient for OLTP workloads involving many small writes, where contention can be an issue when a single page contains many rows. Smaller pages might also be efficient with SSD storage devices, which typically use small block sizes. Keeping the `InnoDB` page size close to the storage device block size minimizes the amount of unchanged data that is rewritten to disk.

InnoDB Memory Configuration

MySQL allocates memory to various caches and buffers to improve performance of database operations. When allocating memory for `InnoDB`, always consider memory required by the operating system, memory allocated to other applications, and memory allocated for other MySQL buffers and caches. For example, if you use `MyISAM` tables, consider the amount of memory allocated for the key buffer (`key_buffer_size`). For an overview of MySQL buffers and caches, see Section 8.12.3.1, “How MySQL Uses Memory”.

Buffers specific to `InnoDB` are configured using the following parameters:

- `innodb_buffer_pool_size` defines size of the buffer pool, which is the memory area that holds cached data for `InnoDB` tables, indexes, and other auxiliary buffers. The size of the buffer pool is important for system performance, and it is typically recommended that `innodb_buffer_pool_size` is configured to 50 to 75 percent of system memory. The default buffer pool size is 128MB. For additional guidance, see Section 8.12.3.1, “How MySQL Uses Memory”. For information about how to configure `InnoDB` buffer pool size, see Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”. Buffer pool size can be configured at startup or dynamically.

  On systems with a large amount of memory, you can improve concurrency by dividing the buffer pool into multiple buffer pool instances. The number of buffer pool instances is controlled by the by `innodb_buffer_pool_instances` option. By default, `InnoDB` creates one buffer pool instance. The number of buffer pool instances can be configured at startup. For more information, see Section 15.6.3.3, “Configuring Multiple Buffer Pool Instances”.

- `innodb_log_buffer_size` defines the size in bytes of the buffer that `InnoDB` uses to write to the log files on disk. The default size is 16MB. A large log buffer enables large transactions to run without a need to write the log to disk before the transactions commit. If you have transactions that update, insert, or delete many rows, you might consider increasing the size of the log buffer to save disk I/O. `innodb_log_buffer_size` can be configured at startup. For related information, see Section 8.5.4, “Optimizing InnoDB Redo Logging”.

Warning

On 32-bit GNU/Linux x86, be careful not to set memory usage too high. `glibc` may permit the process heap to grow over thread stacks, which crashes your server. It is a risk if the memory allocated to the `mysqld` process for global and per-thread buffers and caches is close to or exceeds 2GB.

A formula similar to the following that calculates global and per-thread memory allocation for MySQL can be used to estimate MySQL memory usage. You may need to modify the formula to account for buffers and caches in your MySQL version and configuration. For an overview of MySQL buffers and caches, see Section 8.12.3.1, “How MySQL Uses Memory”.

```
innodb_buffer_pool_size
+ key_buffer_size
+ max_connections*(sort_buffer_size+read_buffer_size+binlog_cache_size)
+ max_connections*2MB
```

Each thread uses a stack (often 2MB, but only 256KB in MySQL binaries provided by Oracle Corporation.) and in the worst case also uses `sort_buffer_size + read_buffer_size` additional memory.

On Linux, if the kernel is enabled for large page support, `InnoDB` can use large pages to allocate memory for its buffer pool. See Section 8.12.3.2, “Enabling Large Page Support”.

15.6.2 Configuring InnoDB for Read-Only Operation

You can now query `InnoDB` tables where the MySQL data directory is on read-only media, by enabling the `--innodb-read-only` configuration option at server startup.

How to Enable

To prepare an instance for read-only operation, make sure all the necessary information is flushed to the data files before storing it on the read-only medium. Run the server with change buffering disabled (`innodb_change_buffering=0`) and do a slow shutdown.

To enable read-only mode for an entire MySQL instance, specify the following configuration options at server startup:

- `--innodb-read-only=1`
- If the instance is on read-only media such as a DVD or CD, or the `/var` directory is not writeable by all: `--pid-file=`*`path_on_writeable_media`* and `--event-scheduler=disabled`

As of MySQL 8.0, enabling `innodb_read_only` prevents table creation and drop operations for all storage engines. These operations modify data dictionary tables in the `mysql` system database, but those tables use the `InnoDB` storage engine and cannot be modified when `innodb_read_only` is enabled. The same restriction applies to any operation that modifies data dictionary tables, such as `ANALYZE TABLE` and `ALTER TABLE` *`tbl_name`* `ENGINE=`*`engine_name`*.

In addition, other tables in the `mysql` system database use the `InnoDB` storage engine in MySQL 8.0. Making those tables read only results in restrictions on operations that modify them. For example, `CREATE USER`, `GRANT`, `REVOKE`, and `INSTALL PLUGIN` operations are not permitted in read-only mode.

Usage Scenarios

This mode of operation is appropriate in situations such as:

- Distributing a MySQL application, or a set of MySQL data, on a read-only storage medium such as a DVD or CD.
- Multiple MySQL instances querying the same data directory simultaneously, typically in a data warehousing configuration. You might use this technique to avoid bottlenecks that can occur with a heavily loaded MySQL instance, or you might use different configuration options for the various instances to tune each one for particular kinds of queries.
- Querying data that has been put into a read-only state for security or data integrity reasons, such as archived backup data.

Note

This feature is mainly intended for flexibility in distribution and deployment, rather than raw performance based on the read-only aspect. See Section 8.5.3, “Optimizing InnoDB Read-Only Transactions” for ways to tune the performance of read-only queries, which do not require making the entire server read-only.

How It Works

When the server is run in read-only mode through the `--innodb-read-only` option, certain `InnoDB` features and components are reduced or turned off entirely:

- No change buffering is done, in particular no merges from the change buffer. To make sure the change buffer is empty when you prepare the instance for read-only operation, disable change buffering (`innodb_change_buffering=0`) and do a slow shutdown first.
- There is no crash recovery phase at startup. The instance must have performed a slow shutdown before being put into the read-only state.
- Because the redo log is not used in read-only operation, you can set `innodb_log_file_size` to the smallest size possible (1 MB) before making the instance read-only.
- All background threads other than I/O read threads are turned off. As a consequence, a read-only instance cannot encounter any deadlock.
- Information about deadlocks, monitor output, and so on is not written to temporary files. As a consequence, `SHOW ENGINE INNODB STATUS` does not produce any output.
- Changes to configuration option settings that would normally change the behavior of write operations, have no effect when the server is in read-only mode.
- The MVCC processing to enforce isolation levels is turned off. All queries read the latest version of a record, because update and deletes are not possible.
- The undo log is not used. Disable any settings for the `innodb_undo_tablespaces` and `innodb_undo_directory` configuration options.

15.6.3 InnoDB Buffer Pool Configuration

This section provides configuration and tuning information for the `InnoDB` buffer pool.

15.6.3.1 The InnoDB Buffer Pool

`InnoDB` maintains a storage area called the buffer pool for caching data and indexes in memory. Knowing how the `InnoDB` buffer pool works, and taking advantage of it to keep frequently accessed data in memory, is an important aspect of MySQL tuning. For information about how the `InnoDB` buffer pool works, see InnoDB Buffer Pool LRU Algorithm.

You can configure the various aspects of the `InnoDB` buffer pool to improve performance.

- Ideally, you set the size of the buffer pool to as large a value as practical, leaving enough memory for other processes on the server to run without excessive paging. The larger the buffer pool, the more `InnoDB` acts like an in-memory database, reading data from disk once and then accessing the data from memory during subsequent reads. See Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”.
- With 64-bit systems with large memory sizes, you can split the buffer pool into multiple parts, to minimize contention for the memory structures among concurrent operations. For details, see Section 15.6.3.3, “Configuring Multiple Buffer Pool Instances”.
- You can keep frequently accessed data in memory despite sudden spikes of activity for operations such as backups or reporting. For details, see Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”.
- You can control when and how `InnoDB` performs read-ahead requests to prefetch pages into the buffer pool asynchronously, in anticipation that the pages will be needed soon. For details, see Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”.
- You can control when background flushing of dirty pages occurs and whether or not `InnoDB` dynamically adjusts the rate of flushing based on workload. For details, see Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing”.
- You can fine-tune aspects of `InnoDB` buffer pool flushing behavior to improve performance. For details, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- You can configure how InnoDB preserves the current buffer pool state to avoid a lengthy warmup period after a server restart. You can also save the current buffer pool state while the server is running. For details, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

InnoDB Buffer Pool LRU Algorithm

InnoDB manages the buffer pool as a list, using a variation of the least recently used (LRU) algorithm. When room is needed to add a new page to the pool, InnoDB evicts the least recently used page and adds the new page to the middle of the list. This “midpoint insertion strategy” treats the list as two sublists:

- At the head, a sublist of “new” (or “young”) pages that were accessed recently.
- At the tail, a sublist of “old” pages that were accessed less recently.

This algorithm keeps pages that are heavily used by queries in the new sublist. The old sublist contains less-used pages; these pages are candidates for eviction.

The LRU algorithm operates as follows by default:

- 3/8 of the buffer pool is devoted to the old sublist.
- The midpoint of the list is the boundary where the tail of the new sublist meets the head of the old sublist.
- When InnoDB reads a page into the buffer pool, it initially inserts it at the midpoint (the head of the old sublist). A page can be read in because it is required for a user-specified operation such as an SQL query, or as part of a read-ahead operation performed automatically by InnoDB.
- Accessing a page in the old sublist makes it “young”, moving it to the head of the buffer pool (the head of the new sublist). If the page was read in because it was required, the first access occurs immediately and the page is made young. If the page was read in due to read-ahead, the first access does not occur immediately (and might not occur at all before the page is evicted).
- As the database operates, pages in the buffer pool that are not accessed “age” by moving toward the tail of the list. Pages in both the new and old sublists age as other pages are made new. Pages in the old sublist also age as pages are inserted at the midpoint. Eventually, a page that remains unused for long enough reaches the tail of the old sublist and is evicted.

By default, pages read by queries immediately move into the new sublist, meaning they stay in the buffer pool longer. A table scan (such as performed for a mysqldump operation, or a SELECT statement with no WHERE clause) can bring a large amount of data into the buffer pool and evict an equivalent amount of older data, even if the new data is never used again. Similarly, pages that are loaded by the read-ahead background thread and then accessed only once move to the head of the new list. These situations can push frequently used pages to the old sublist, where they become subject to eviction. For information about optimizing this behavior, see Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”, and Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”.

InnoDB Standard Monitor output contains several fields in the BUFFER POOL AND MEMORY section that pertain to operation of the buffer pool LRU algorithm. For details, see Section 15.6.3.9, “Monitoring the Buffer Pool Using the InnoDB Standard Monitor”.

InnoDB Buffer Pool Configuration Options

Several configuration options affect different aspects of the InnoDB buffer pool.

- innodb_buffer_pool_size

  Specifies the size of the buffer pool. If the buffer pool is small and you have sufficient memory, making the buffer pool larger can improve performance by reducing the amount of disk I/O needed as queries access InnoDB tables. The innodb_buffer_pool_size option is dynamic, which

allows you to configure buffer pool size without restarting the server. See Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size” for more information.

- `innodb_buffer_pool_chunk_size`

  Defines the chunk size for `InnoDB` buffer pool resizing operations. See Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size” for more information.

- `innodb_buffer_pool_instances`

  Divides the buffer pool into a user-specified number of separate regions, each with its own LRU list and related data structures, to reduce contention during concurrent memory read and write operations. This option only takes effect when you set `innodb_buffer_pool_size` to a value of 1GB or more. The total size you specify is divided among all the buffer pools. For best efficiency, specify a combination of `innodb_buffer_pool_instances` and `innodb_buffer_pool_size` so that each buffer pool instance is at least 1 gigabyte. See Section 15.6.3.3, “Configuring Multiple Buffer Pool Instances” for more information.

- `innodb_old_blocks_pct`

  Specifies the approximate percentage of the buffer pool that `InnoDB` uses for the old block sublist. The range of values is 5 to 95. The default value is 37 (that is, 3/8 of the pool). See Section 15.6.3.4, “Making the Buffer Pool Scan Resistant” for more information.

- `innodb_old_blocks_time`

  Specifies how long in milliseconds (ms) a page inserted into the old sublist must stay there after its first access before it can be moved to the new sublist. If the value is 0, a page inserted into the old sublist moves immediately to the new sublist the first time it is accessed, no matter how soon after insertion the access occurs. If the value is greater than 0, pages remain in the old sublist until an access occurs at least that many milliseconds after the first access. For example, a value of 1000 causes pages to stay in the old sublist for 1 second after the first access before they become eligible to move to the new sublist.

  Setting `innodb_old_blocks_time` greater than 0 prevents one-time table scans from flooding the new sublist with pages used only for the scan. Rows in a page read in for a scan are accessed many times in rapid succession, but the page is unused after that. If `innodb_old_blocks_time` is set to a value greater than time to process the page, the page remains in the “old” sublist and ages to the tail of the list to be evicted quickly. This way, pages used only for a one-time scan do not act to the detriment of heavily used pages in the new sublist.

  `innodb_old_blocks_time` can be set at runtime, so you can change it temporarily while performing operations such as table scans and dumps:

```
SET GLOBAL innodb_old_blocks_time = 1000;
... perform queries that scan tables ...
SET GLOBAL innodb_old_blocks_time = 0;
```

  This strategy does not apply if your intent is to “warm up” the buffer pool by filling it with a table's content. For example, benchmark tests often perform a table or index scan at server startup, because that data would normally be in the buffer pool after a period of normal use. In this case, leave `innodb_old_blocks_time` set to 0, at least until the warmup phase is complete.

  See Section 15.6.3.4, “Making the Buffer Pool Scan Resistant” for more information.

- `innodb_read_ahead_threshold`

  Controls the sensitivity of linear read-ahead that `InnoDB` uses to prefetch pages into the buffer pool.

  See Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)” for more information.

- `innodb_random_read_ahead`

  Enables random read-ahead technique for prefetching pages into the buffer pool. Random read-ahead is a technique that predicts when pages might be needed soon based on pages already in the buffer pool, regardless of the order in which those pages were read. `innodb_random_read_ahead` is disabled by default.

  See Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)” for more information.

- `innodb_adaptive_flushing`

  Specifies whether to dynamically adjust the rate of flushing dirty pages in the buffer pool based on workload. Adjusting the flush rate dynamically is intended to avoid bursts of I/O activity. This setting is enabled by default.

  See Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing” for more information.

- `innodb_adaptive_flushing_lwm`

  Low water mark representing percentage of redo log capacity at which adaptive flushing is enabled.

  See Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing” for more information.

- `innodb_flush_neighbors`

  Specifies whether flushing a page from the buffer pool also flushes other dirty pages in the same extent.

  See Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing” for more information.

- `innodb_flushing_avg_loops`

  Number of iterations for which InnoDB keeps the previously calculated snapshot of the flushing state, controlling how quickly adaptive flushing responds to changing workloads.

  See Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing” for more information.

- `innodb_lru_scan_depth`

  A parameter that influences the algorithms and heuristics for the flush operation for the buffer pool. Primarily of interest to performance experts tuning I/O-intensive workloads. It specifies, per buffer pool instance, how far down the buffer pool LRU list the `page_cleaner` thread scans looking for dirty pages to flush.

  See Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing” for more information.

- `innodb_max_dirty_pages_pct`

  `InnoDB` tries to flush data from the buffer pool so that the percentage of dirty pages does not exceed this value. Specify an integer in the range from 0 to 99. The default value is 75.

  See Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing” for more information.

- `innodb_max_dirty_pages_pct_lwm`

  Low water mark representing percentage of dirty pages where preflushing is enabled to control the dirty page ratio. The default of 0 disables the pre-flushing behavior entirely.

  See Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing” for more information.

- `innodb_buffer_pool_filename`

Specifies the name of the file that holds the list of tablespace IDs and page IDs produced by `innodb_buffer_pool_dump_at_shutdown` or `innodb_buffer_pool_dump_now`.

See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_dump_at_shutdown`

  Specifies whether to record the pages cached in the buffer pool when the MySQL server is shut down, to shorten the warmup process at the next restart.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_load_at_startup`

  Specifies that, on MySQL server startup, the buffer pool is automatically warmed up by loading the same pages it held at an earlier time. Typically used in combination with `innodb_buffer_pool_dump_at_shutdown`.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_dump_now`

  Immediately records the pages cached in the buffer pool.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_load_now`

  Immediately warms up the buffer pool by loading a set of data pages, without waiting for a server restart. Can be useful to bring cache memory back to a known state during benchmarking, or to ready the MySQL server to resume its normal workload after running queries for reports or maintenance. Typically used with `innodb_buffer_pool_dump_now`.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_dump_pct`

  Specifies the percentage of the most recently used pages for each buffer pool to read out and dump. The range is 1 to 100.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

- `innodb_buffer_pool_load_abort`

  Interrupts the process of restoring buffer pool contents triggered by `innodb_buffer_pool_load_at_startup` or `innodb_buffer_pool_load_now`.

  See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State” for more information.

15.6.3.2 Configuring InnoDB Buffer Pool Size

You can configure `InnoDB` buffer pool size offline (at startup) or online, while the server is running. Behavior described in this section applies to both methods. For additional information about configuring buffer pool size online, see Configuring InnoDB Buffer Pool Size Online.

When increasing or decreasing `innodb_buffer_pool_size`, the operation is performed in chunks. Chunk size is defined by the `innodb_buffer_pool_chunk_size` configuration option, which has a default of `128M`. For more information, see Configuring InnoDB Buffer Pool Chunk Size.

Buffer pool size must always be equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances`. If you configure `innodb_buffer_pool_size` to a value that is not equal to or a multiple of `innodb_buffer_pool_chunk_size` *

`innodb_buffer_pool_instances`, buffer pool size is automatically adjusted to a value that is equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` that is not less than the specified buffer pool size.

In the following example, `innodb_buffer_pool_size` is set to `8G`, and `innodb_buffer_pool_instances` is set to `16`. `innodb_buffer_pool_chunk_size` is `128M`, which is the default value.

`8G` is a valid `innodb_buffer_pool_size` value because `8G` is a multiple of `innodb_buffer_pool_instances=16` * `innodb_buffer_pool_chunk_size=128M`, which is `2G`.

```
shell> mysqld --innodb_buffer_pool_size=8G --innodb_buffer_pool_instances=16

mysql> SELECT @@innodb_buffer_pool_size/1024/1024/1024;
+----------------------------------------+
| @@innodb_buffer_pool_size/1024/1024/1024 |
+----------------------------------------+
|                          8.000000000000 |
+----------------------------------------+
```

In this example, `innodb_buffer_pool_size` is set to `9G`, and `innodb_buffer_pool_instances` is set to `16`. `innodb_buffer_pool_chunk_size` is `128M`, which is the default value. In this case, `9G` is not a multiple of `innodb_buffer_pool_instances=16` * `innodb_buffer_pool_chunk_size=128M`, so `innodb_buffer_pool_size` is adjusted to `10G`, which is the next multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` that is not less than the specified buffer pool size.

```
shell> mysqld --innodb_buffer_pool_size=9G --innodb_buffer_pool_instances=16

mysql> SELECT @@innodb_buffer_pool_size/1024/1024/1024;
+----------------------------------------+
| @@innodb_buffer_pool_size/1024/1024/1024 |
+----------------------------------------+
|                         10.000000000000 |
+----------------------------------------+
```

Configuring InnoDB Buffer Pool Chunk Size

`innodb_buffer_pool_chunk_size` can be increased or decreased in 1MB (1048576 byte) units but can only be modified at startup, in a command line string or in a MySQL configuration file.

Command line:

```
shell> mysqld --innodb_buffer_pool_chunk_size=134217728
```

Configuration file:

```
[mysqld]
innodb_buffer_pool_chunk_size=134217728
```

The following conditions apply when altering `innodb_buffer_pool_chunk_size`:

- If the new `innodb_buffer_pool_chunk_size` value * `innodb_buffer_pool_instances` is larger than the current buffer pool size when the buffer pool is initialized, `innodb_buffer_pool_chunk_size` is truncated to `innodb_buffer_pool_size` / `innodb_buffer_pool_instances`.

  For example, if the buffer pool is initialized with a size of `2GB` (2147483648 bytes), `4` buffer pool instances, and a chunk size of `1GB` (1073741824 bytes), chunk size is truncated to a value equal to `innodb_buffer_pool_size` / `innodb_buffer_pool_instances`, as shown below:

  ```
  shell> mysqld --innodb_buffer_pool_size=2147483648 --innodb_buffer_pool_instances=4
  --innodb_buffer_pool_chunk_size=1073741824;

  mysql> SELECT @@innodb_buffer_pool_size;
  +---------------------------+
  ```

```
| @@innodb_buffer_pool_size |
+---------------------------+
|                2147483648 |
+---------------------------+

mysql> SELECT @@innodb_buffer_pool_instances;
+---------------------------------+
| @@innodb_buffer_pool_instances |
+---------------------------------+
|                               4 |
+---------------------------------+

# Chunk size was set to 1GB (1073741824 bytes) on startup but was
# truncated to innodb_buffer_pool_size / innodb_buffer_pool_instances

mysql> SELECT @@innodb_buffer_pool_chunk_size;
+---------------------------------+
| @@innodb_buffer_pool_chunk_size |
+---------------------------------+
|                       536870912 |
+---------------------------------+
```

- Buffer pool size must always be equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances`. If you alter `innodb_buffer_pool_chunk_size`, `innodb_buffer_pool_size` is automatically adjusted to a value that is equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` that is not less than current buffer pool size. The adjustment occurs when the buffer pool is initialized. This behavior is demonstrated in the following example:

```
# The buffer pool has a default size of 128MB (134217728 bytes)

mysql> SELECT @@innodb_buffer_pool_size;
+---------------------------+
| @@innodb_buffer_pool_size |
+---------------------------+
|                 134217728 |
+---------------------------+

# The chunk size is also 128MB (134217728 bytes)

mysql> SELECT @@innodb_buffer_pool_chunk_size;
+---------------------------------+
| @@innodb_buffer_pool_chunk_size |
+---------------------------------+
|                       134217728 |
+---------------------------------+

# There is a single buffer pool instance

mysql> SELECT @@innodb_buffer_pool_instances;
+---------------------------------+
| @@innodb_buffer_pool_instances |
+---------------------------------+
|                               1 |
+---------------------------------+

# Chunk size is decreased by 1MB (1048576 bytes) at startup
# (134217728 - 1048576 = 133169152):

shell> mysqld --innodb_buffer_pool_chunk_size=133169152

mysql> SELECT @@innodb_buffer_pool_chunk_size;
+---------------------------------+
| @@innodb_buffer_pool_chunk_size |
+---------------------------------+
|                       133169152 |
+---------------------------------+

# Buffer pool size increases from 134217728 to 266338304
# Buffer pool size is automatically adjusted to a value that is equal to
# or a multiple of innodb_buffer_pool_chunk_size * innodb_buffer_pool_instances
```

```
# that is not less than current buffer pool size

mysql> SELECT @@innodb_buffer_pool_size;
+---------------------------+
| @@innodb_buffer_pool_size |
+---------------------------+
|                 266338304 |
+---------------------------+
```

This example demonstrates the same behavior but with multiple buffer pool instances:

```
# The buffer pool has a default size of 2GB (2147483648 bytes)

mysql> SELECT @@innodb_buffer_pool_size;
+---------------------------+
| @@innodb_buffer_pool_size |
+---------------------------+
|                2147483648 |
+---------------------------+

# The chunk size is .5 GB (536870912 bytes)

mysql> SELECT @@innodb_buffer_pool_chunk_size;
+--------------------------------+
| @@innodb_buffer_pool_chunk_size |
+--------------------------------+
|                      536870912 |
+--------------------------------+

# There are 4 buffer pool instances

mysql> SELECT @@innodb_buffer_pool_instances;
+-------------------------------+
| @@innodb_buffer_pool_instances |
+-------------------------------+
|                             4 |
+-------------------------------+

# Chunk size is decreased by 1MB (1048576 bytes) at startup
# (536870912 - 1048576 = 535822336):

shell> mysqld --innodb_buffer_pool_chunk_size=535822336

mysql> SELECT @@innodb_buffer_pool_chunk_size;
+--------------------------------+
| @@innodb_buffer_pool_chunk_size |
+--------------------------------+
|                      535822336 |
+--------------------------------+

# Buffer pool size increases from 2147483648 to 4286578688
# Buffer pool size is automatically adjusted to a value that is equal to
# or a multiple of innodb_buffer_pool_chunk_size * innodb_buffer_pool_instances
# that is not less than current buffer pool size of 2147483648

mysql> SELECT @@innodb_buffer_pool_size;
+---------------------------+
| @@innodb_buffer_pool_size |
+---------------------------+
|                4286578688 |
+---------------------------+
```

Care should be taken when changing `innodb_buffer_pool_chunk_size`, as changing this value can increase the size of the buffer pool, as shown in the examples above. Before you change `innodb_buffer_pool_chunk_size`, calculate the effect on `innodb_buffer_pool_size` to ensure that the resulting buffer pool size is acceptable.

Note

To avoid potential performance issues, the number of chunks (`innodb_buffer_pool_size` / `innodb_buffer_pool_chunk_size`) should not exceed 1000.

Configuring InnoDB Buffer Pool Size Online

The `innodb_buffer_pool_size` configuration option can be set dynamically using a `SET` statement, allowing you to resize the buffer pool without restarting the server. For example:

```
mysql> SET GLOBAL innodb_buffer_pool_size=402653184;
```

Active transactions and operations performed through `InnoDB` APIs should be completed before resizing the buffer pool. When initiating a resizing operation, the operation does not start until all active transactions are completed. Once the resizing operation is in progress, new transactions and operations that require access to the buffer pool must wait until the resizing operation finishes. The exception to the rule is that concurrent access to the buffer pool is permitted while the buffer pool is defragmented and pages are withdrawn when buffer pool size is decreased. A drawback of allowing concurrent access is that it could result in a temporary shortage of available pages while pages are being withdrawn.

Note

Nested transactions could fail if initiated after the buffer pool resizing operation begins.

Monitoring Online Buffer Pool Resizing Progress

The `Innodb_buffer_pool_resize_status` reports buffer pool resizing progress. For example:

```
mysql> SHOW STATUS WHERE Variable_name='InnoDB_buffer_pool_resize_status';
+----------------------------------+----------------------------------+
| Variable_name                    | Value                            |
+----------------------------------+----------------------------------+
| Innodb_buffer_pool_resize_status | Resizing also other hash tables. |
+----------------------------------+----------------------------------+
```

Buffer pool resizing progress is also logged in the server error log file. This example shows notes that are logged when increasing the size of the buffer pool:

```
[Note] InnoDB: Resizing buffer pool from 134217728 to 4294967296. (unit=134217728)
[Note] InnoDB: disabled adaptive hash index.
[Note] InnoDB: buffer pool 0 : 31 chunks (253952 blocks) was added.
[Note] InnoDB: buffer pool 0 : hash tables were resized.
[Note] InnoDB: Resized hash tables at lock_sys, adaptive hash index, dictionary.
[Note] InnoDB: completed to resize buffer pool from 134217728 to 4294967296.
[Note] InnoDB: re-enabled adaptive hash index.
```

This example shows notes that are logged when decreasing the size of the buffer pool:

```
[Note] InnoDB: Resizing buffer pool from 4294967296 to 134217728. (unit=134217728)
[Note] InnoDB: disabled adaptive hash index.
[Note] InnoDB: buffer pool 0 : start to withdraw the last 253952 blocks.
[Note] InnoDB: buffer pool 0 : withdrew 253952 blocks from free list. tried to relocate 0 pages.
(253952/253952)
[Note] InnoDB: buffer pool 0 : withdrawn target 253952 blocks.
[Note] InnoDB: buffer pool 0 : 31 chunks (253952 blocks) was freed.
[Note] InnoDB: buffer pool 0 : hash tables were resized.
[Note] InnoDB: Resized hash tables at lock_sys, adaptive hash index, dictionary.
[Note] InnoDB: completed to resize buffer pool from 4294967296 to 134217728.
[Note] InnoDB: re-enabled adaptive hash index.
```

Online Buffer Pool Resizing Internals

The resizing operation is performed by a background thread. When increasing the size of the buffer pool, the resizing operation:

- Adds pages in `chunks` (chunk size is defined by `innodb_buffer_pool_chunk_size`)
- Coverts hash tables, lists, and pointers to use new addresses in memory
- Adds new pages to the free list

While these operations are in progress, other threads are blocked from accessing the buffer pool.

When decreasing the size of the buffer pool, the resizing operation:

- Defragments the buffer pool and withdraws (frees) pages
- Removes pages in `chunks` (chunk size is defined by `innodb_buffer_pool_chunk_size`)
- Converts hash tables, lists, and pointers to use new addresses in memory

Of these operations, only defragmenting the buffer pool and withdrawing pages allow other threads to access to the buffer pool concurrently.

15.6.3.3 Configuring Multiple Buffer Pool Instances

For systems with buffer pools in the multi-gigabyte range, dividing the buffer pool into separate instances can improve concurrency, by reducing contention as different threads read and write to cached pages. This feature is typically intended for systems with a buffer pool size in the multi-gigabyte range. Multiple buffer pool instances are configured using the `innodb_buffer_pool_instances` configuration option, and you might also adjust the `innodb_buffer_pool_size` value.

When the `InnoDB` buffer pool is large, many data requests can be satisfied by retrieving from memory. You might encounter bottlenecks from multiple threads trying to access the buffer pool at once. You can enable multiple buffer pools to minimize this contention. Each page that is stored in or read from the buffer pool is assigned to one of the buffer pools randomly, using a hashing function. Each buffer pool manages its own free lists, flush lists, LRUs, and all other data structures connected to a buffer pool. Prior to MySQL 8.0, each buffer pool was protected by its own buffer pool mutex. In MySQL 8.0 and later, the buffer pool mutex was replaced by several list and hash protecting mutexes, to reduce contention.

To enable multiple buffer pool instances, set the `innodb_buffer_pool_instances` configuration option to a value greater than 1 (the default) up to 64 (the maximum). This option takes effect only when you set `innodb_buffer_pool_size` to a size of 1GB or more. The total size you specify is divided among all the buffer pools. For best efficiency, specify a combination of `innodb_buffer_pool_instances` and `innodb_buffer_pool_size` so that each buffer pool instance is at least 1GB.

For information about modifying `InnoDB` buffer pool size, see Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”.

15.6.3.4 Making the Buffer Pool Scan Resistant

Rather than using a strict LRU algorithm, `InnoDB` uses a technique to minimize the amount of data that is brought into the buffer pool and never accessed again. The goal is to make sure that frequently accessed (“hot”) pages remain in the buffer pool, even as read-ahead and full table scans bring in new blocks that might or might not be accessed afterward.

Newly read blocks are inserted into the middle of the LRU list. All newly read pages are inserted at a location that by default is `3/8` from the tail of the LRU list. The pages are moved to the front of the list (the most-recently used end) when they are accessed in the buffer pool for the first time. Thus, pages that are never accessed never make it to the front portion of the LRU list, and “age out” sooner than with a strict LRU approach. This arrangement divides the LRU list into two segments, where the pages downstream of the insertion point are considered “old” and are desirable victims for LRU eviction.

For an explanation of the inner workings of the InnoDB buffer pool and specifics about the LRU algorithm, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

You can control the insertion point in the LRU list and choose whether InnoDB applies the same optimization to blocks brought into the buffer pool by table or index scans. The configuration parameter innodb_old_blocks_pct controls the percentage of “old” blocks in the LRU list. The default value of innodb_old_blocks_pct is 37, corresponding to the original fixed ratio of 3/8. The value range is 5 (new pages in the buffer pool age out very quickly) to 95 (only 5% of the buffer pool is reserved for hot pages, making the algorithm close to the familiar LRU strategy).

The optimization that keeps the buffer pool from being churned by read-ahead can avoid similar problems due to table or index scans. In these scans, a data page is typically accessed a few times in quick succession and is never touched again. The configuration parameter innodb_old_blocks_time specifies the time window (in milliseconds) after the first access to a page during which it can be accessed without being moved to the front (most-recently used end) of the LRU list. The default value of innodb_old_blocks_time is 1000. Increasing this value makes more and more blocks likely to age out faster from the buffer pool.

Both innodb_old_blocks_pct and innodb_old_blocks_time are dynamic, global and can be specified in the MySQL option file (my.cnf or my.ini) or changed at runtime with the SET GLOBAL command. Changing the setting requires the SYSTEM_VARIABLES_ADMIN or SUPER privilege.

To help you gauge the effect of setting these parameters, the SHOW ENGINE INNODB STATUS command reports buffer pool statistics. For details, see Section 15.6.3.9, “Monitoring the Buffer Pool Using the InnoDB Standard Monitor”.

Because the effects of these parameters can vary widely based on your hardware configuration, your data, and the details of your workload, always benchmark to verify the effectiveness before changing these settings in any performance-critical or production environment.

In mixed workloads where most of the activity is OLTP type with periodic batch reporting queries which result in large scans, setting the value of innodb_old_blocks_time during the batch runs can help keep the working set of the normal workload in the buffer pool.

When scanning large tables that cannot fit entirely in the buffer pool, setting innodb_old_blocks_pct to a small value keeps the data that is only read once from consuming a significant portion of the buffer pool. For example, setting innodb_old_blocks_pct=5 restricts this data that is only read once to 5% of the buffer pool.

When scanning small tables that do fit into memory, there is less overhead for moving pages around within the buffer pool, so you can leave innodb_old_blocks_pct at its default value, or even higher, such as innodb_old_blocks_pct=50.

The effect of the innodb_old_blocks_time parameter is harder to predict than the innodb_old_blocks_pct parameter, is relatively small, and varies more with the workload. To arrive at an optimal value, conduct your own benchmarks if the performance improvement from adjusting innodb_old_blocks_pct is not sufficient.

15.6.3.5 Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)

A read-ahead request is an I/O request to prefetch multiple pages in the buffer pool asynchronously, in anticipation that these pages will be needed soon. The requests bring in all the pages in one extent. InnoDB uses two read-ahead algorithms to improve I/O performance:

Linear read-ahead is a technique that predicts what pages might be needed soon based on pages in the buffer pool being accessed sequentially. You control when InnoDB performs a read-ahead operation by adjusting the number of sequential page accesses required to trigger an asynchronous read request, using the configuration parameter innodb_read_ahead_threshold. Before this

parameter was added, InnoDB would only calculate whether to issue an asynchronous prefetch request for the entire next extent when it read in the last page of the current extent.

The configuration parameter innodb_read_ahead_threshold controls how sensitive InnoDB is in detecting patterns of sequential page access. If the number of pages read sequentially from an extent is greater than or equal to innodb_read_ahead_threshold, InnoDB initiates an asynchronous read-ahead operation of the entire following extent. innodb_read_ahead_threshold can be set to any value from 0-64. The default value is 56. The higher the value, the more strict the access pattern check. For example, if you set the value to 48, InnoDB triggers a linear read-ahead request only when 48 pages in the current extent have been accessed sequentially. If the value is 8, InnoDB triggers an asynchronous read-ahead even if as few as 8 pages in the extent are accessed sequentially. You can set the value of this parameter in the MySQL configuration file, or change it dynamically with the SET GLOBAL command, which requires the SYSTEM_VARIABLES_ADMIN or SUPER privilege.

Random read-ahead is a technique that predicts when pages might be needed soon based on pages already in the buffer pool, regardless of the order in which those pages were read. If 13 consecutive pages from the same extent are found in the buffer pool, InnoDB asynchronously issues a request to prefetch the remaining pages of the extent. To enable this feature, set the configuration variable innodb_random_read_ahead to ON.

The SHOW ENGINE INNODB STATUS command displays statistics to help you evaluate the effectiveness of the read-ahead algorithm. Statistics include counter information for the following global status variables:

- Innodb_buffer_pool_read_ahead
- Innodb_buffer_pool_read_ahead_evicted
- Innodb_buffer_pool_read_ahead_rnd

This information can be useful when fine-tuning the innodb_random_read_ahead setting.

For more information about I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O” and Section 8.12.1, “Optimizing Disk I/O”.

15.6.3.6 Configuring InnoDB Buffer Pool Flushing

InnoDB performs certain tasks in the background, including flushing of dirty pages (those pages that have been changed but are not yet written to the database files) from the buffer pool.

InnoDB starts flushing buffer pool pages when the percentage of dirty pages in the buffer pool reaches the low water mark setting defined by innodb_max_dirty_pages_pct_lwm. This option is intended to control the ratio of dirty pages in the buffer pool and ideally prevent the percentage of dirty pages from reaching innodb_max_dirty_pages_pct. If the percentage of dirty pages in the buffer pool exceeds innodb_max_dirty_pages_pct, InnoDB begins to aggressively flush buffer pool pages.

InnoDB uses an algorithm to estimate the required rate of flushing, based on the speed of redo log generation and the current rate of flushing. The intent is to smooth overall performance by ensuring that buffer flush activity keeps up with the need to keep the buffer pool “clean”. Automatically adjusting the rate of flushing can help to avoid sudden dips in throughput, when excessive buffer pool flushing limits the I/O capacity available for ordinary read and write activity.

InnoDB uses its log files in a circular fashion. Before reusing a portion of a log file, InnoDB flushes to disk all dirty buffer pool pages whose redo entries are contained in that portion of the log file, a process known as a sharp checkpoint. If a workload is write-intensive, it generates a lot of redo information, all written to the log file. If all available space in the log files is used up, a sharp checkpoint occurs, causing a temporary reduction in throughput. This situation can happen even if innodb_max_dirty_pages_pct is not reached.

InnoDB uses a heuristic-based algorithm to avoid such a scenario, by measuring the number of dirty pages in the buffer pool and the rate at which redo is being generated. Based on these numbers,

InnoDB decides how many dirty pages to flush from the buffer pool each second. This self-adapting algorithm is able to deal with sudden changes in workload.

Internal benchmarking has shown that this algorithm not only maintains throughput over time, but can also improve overall throughput significantly.

Because adaptive flushing can significantly affect the I/O pattern of a workload, the innodb_adaptive_flushing configuration parameter lets you turn off this feature. The default value for innodb_adaptive_flushing is ON, enabling the adaptive flushing algorithm. You can set the value of this parameter in the MySQL option file (my.cnf or my.ini) or change it dynamically with the SET GLOBAL command, which requires the SYSTEM_VARIABLES_ADMIN or SUPER privilege.

For information about fine-tuning InnoDB buffer pool flushing behavior, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

For more information about InnoDB I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

15.6.3.7 Fine-tuning InnoDB Buffer Pool Flushing

The configuration options innodb_flush_neighbors and innodb_lru_scan_depth let you fine-tune certain aspects of the flushing process for the InnoDB buffer pool. These options primarily help write-intensive workloads. With heavy DML activity, flushing can fall behind if it is not aggressive enough, resulting in excessive memory use in the buffer pool; or, disk writes due to flushing can saturate your I/O capacity if that mechanism is too aggressive. The ideal settings depend on your workload, data access patterns, and storage configuration (for example, whether data is stored on HDD or SSD devices).

For systems with constant heavy workloads, or workloads that fluctuate widely, several configuration options let you fine-tune the flushing behavior for InnoDB tables:

- innodb_adaptive_flushing_lwm
- innodb_max_dirty_pages_pct_lwm
- innodb_io_capacity_max
- innodb_flushing_avg_loops

These options feed into the formula used by the innodb_adaptive_flushing option.

The innodb_adaptive_flushing, innodb_io_capacity and innodb_max_dirty_pages_pct options are limited or extended by the following options:

- innodb_adaptive_flushing_lwm
- innodb_io_capacity_max
- innodb_max_dirty_pages_pct_lwm

The InnoDB adaptive flushing mechanism is not appropriate in all cases. It gives the most benefit when the redo log is in danger of filling up. The innodb_adaptive_flushing_lwm option specifies a “low water mark” percentage of redo log capacity; when that threshold is crossed, InnoDB turns on adaptive flushing even if not specified by the innodb_adaptive_flushing option.

If flushing activity falls far behind, InnoDB can flush more aggressively than specified by innodb_io_capacity. innodb_io_capacity_max represents an upper limit on the I/O capacity used in such emergency situations, so that the spike in I/O does not consume all the capacity of the server.

InnoDB tries to flush data from the buffer pool so that the percentage of dirty pages does not exceed the value of innodb_max_dirty_pages_pct. The default value for innodb_max_dirty_pages_pct is 75.

Note

The `innodb_max_dirty_pages_pct` setting establishes a target for flushing activity. It does not affect the rate of flushing. For information about managing the rate of flushing, see Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing”.

The `innodb_max_dirty_pages_pct_lwm` option specifies a “low water mark” value that represents the percentage of dirty pages where pre-flushing is enabled to control the dirty page ratio and ideally prevent the percentage of dirty pages from reaching `innodb_max_dirty_pages_pct`. A value of `innodb_max_dirty_pages_pct_lwm=0` disables the “pre-flushing” behavior.

Most of the options referenced above are most applicable to servers that run write-heavy workloads for long periods of time and have little reduced load time to catch up with changes waiting to be written to disk.

`innodb_flushing_avg_loops` defines the number of iterations for which `InnoDB` keeps the previously calculated snapshot of the flushing state, which controls how quickly adaptive flushing responds to foreground load changes. Setting a high value for `innodb_flushing_avg_loops` means that `InnoDB` keeps the previously calculated snapshot longer, so adaptive flushing responds more slowly. A high value also reduces positive feedback between foreground and background work, but when setting a high value it is important to ensure that `InnoDB` redo log utilization does not reach 75% (the hardcoded limit at which async flushing starts) and that the `innodb_max_dirty_pages_pct` setting keeps the number of dirty pages to a level that is appropriate for the workload.

Systems with consistent workloads, a large `innodb_log_file_size`, and small spikes that do not reach 75% redo log space utilization should use a high `innodb_flushing_avg_loops` value to keep flushing as smooth as possible. For systems with extreme load spikes or log files that do not provide a lot of space, consider a smaller `innodb_flushing_avg_loops` value. A smaller value allows flushing to closely track the load and helps avoid reaching 75% redo log space utilization.

15.6.3.8 Saving and Restoring the Buffer Pool State

To reduce the warmup period after restarting the server, `InnoDB` saves a percentage of the most recently used pages for each buffer pool at server shutdown and restores these pages at server startup. The percentage of recently used pages that is stored is defined by the `innodb_buffer_pool_dump_at_shutdown` configuration option.

After restarting a busy server, there is typically a warmup period with steadily increasing throughput, as disk pages that were in the buffer pool are brought back into memory (as the same data is queried, updated, and so on). The ability to restore the buffer pool at startup shortens the warmup period by reloading disk pages that were in the buffer pool before the restart rather than waiting for DML operations to access corresponding rows. Also, I/O requests can be performed in large batches, making the overall I/O faster. Page loading happens in the background, and does not delay database startup.

In addition to saving the buffer pool state at shutdown and restoring it at startup, you can save and restore the buffer pool state at any time, while the server is running. For example, you can save the state of the buffer pool after reaching a stable throughput under a steady workload. You could also restore the previous buffer pool state after running reports or maintenance jobs that bring data pages into the buffer pool that are only requited for those operations, or after running some other non-typical workload.

Even though a buffer pool can be many gigabytes in size, the buffer pool data that `InnoDB` saves to disk is tiny by comparison. Only tablespace IDs and page IDs necessary to locate the appropriate pages are saved to disk. This information is derived from the `INNODB_BUFFER_PAGE_LRU INFORMATION_SCHEMA` table. By default, tablespace ID and page ID data is saved in a file named `ib_buffer_pool`, which is saved to the `InnoDB` data directory. The file name and location can be modified using the `innodb_buffer_pool_filename` configuration parameter.

Because data is cached in and aged out of the buffer pool as it is with regular database operations, there is no problem if the disk pages are recently updated, or if a DML operation involves data that has not yet been loaded. The loading mechanism skips requested pages that no longer exist.

The underlying mechanism involves a background thread that is dispatched to perform the dump and load operations.

Disk pages from compressed tables are loaded into the buffer pool in their compressed form. Pages are uncompressed as usual when page contents are accessed during DML operations. Because uncompressing pages is a CPU-intensive process, it is more efficient for concurrency to perform the operation in a connection thread rather than in the single thread that performs the buffer pool restore operation.

Operations related to saving and restoring the buffer pool state are described in the following topics:

- Configuring the Dump Percentage for Buffer Pool Pages
- Saving the Buffer Pool State at Shutdown and Restoring it at Startup
- Saving and Restoring the Buffer Pool State Online
- Displaying Buffer Pool Dump Progress
- Displaying Buffer Pool Load Progress
- Aborting a Buffer Pool Load Operation
- Monitoring Buffer Pool Load Progress Using Performance Schema

Configuring the Dump Percentage for Buffer Pool Pages

Before dumping pages from the buffer pool, you can configure the percentage of most-recently-used buffer pool pages that you want to dump by setting the `innodb_buffer_pool_dump_pct` option. If you plan to dump buffer pool pages while the server is running, you can configure the option dynamically:

```
SET GLOBAL innodb_buffer_pool_dump_pct=40;
```

If you plan to dump buffer pool pages at server shutdown, set `innodb_buffer_pool_dump_pct` in your configuration file.

```
[mysqld]
      innodb_buffer_pool_dump_pct=40
```

The `innodb_buffer_pool_dump_pct` default value is 25 (dump 25% of most-recently-used pages).

Saving the Buffer Pool State at Shutdown and Restoring it at Startup

To save the state of the buffer pool at server shutdown, issue the following statement prior to shutting down the server:

```
SET GLOBAL innodb_buffer_pool_dump_at_shutdown=ON;
```

`innodb_buffer_pool_dump_at_shutdown` is enabled by default.

To restore the buffer pool state at server startup, specify the `--innodb_buffer_pool_load_at_startup` option when starting the server:

```
mysqld --innodb_buffer_pool_load_at_startup=ON;
```

`innodb_buffer_pool_load_at_startup` is enabled by default.

Saving and Restoring the Buffer Pool State Online

To save the state of the buffer pool while MySQL server is running, issue the following statement:

```
SET GLOBAL innodb_buffer_pool_dump_now=ON;
```

To restore the buffer pool state while MySQL is running, issue the following statement:

```
SET GLOBAL innodb_buffer_pool_load_now=ON;
```

Displaying Buffer Pool Dump Progress

To display progress when saving the buffer pool state to disk, issue the following statement:

```
SHOW STATUS LIKE 'Innodb_buffer_pool_dump_status';
```

If the operation has not yet started, “not started” is returned. If the operation is complete, the completion time is printed (e.g. Finished at 110505 12:18:02). If the operation is in progress, status information is provided (e.g. Dumping buffer pool 5/7, page 237/2873).

Displaying Buffer Pool Load Progress

To display progress when loading the buffer pool, issue the following statement:

```
SHOW STATUS LIKE 'Innodb_buffer_pool_load_status';
```

If the operation has not yet started, “not started” is returned. If the operation is complete, the completion time is printed (e.g. Finished at 110505 12:23:24). If the operation is in progress, status information is provided (e.g. Loaded 123/22301 pages).

Aborting a Buffer Pool Load Operation

To abort a buffer pool load operation, issue the following statement:

```
SET GLOBAL innodb_buffer_pool_load_abort=ON;
```

Monitoring Buffer Pool Load Progress Using Performance Schema

You can monitor buffer pool load progress using Performance Schema.

The following example demonstrates how to enable the `stage/innodb/buffer pool load` stage event instrument and related consumer tables to monitor buffer pool load progress.

For information about buffer pool dump and load procedures used in this example, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”. For information about Performance Schema stage event instruments and related consumers, see Section 25.11.5, “Performance Schema Stage Event Tables”.

1. Enable the `stage/innodb/buffer pool load` instrument:

   ```
   mysql> UPDATE performance_schema.setup_instruments SET ENABLED = 'YES'
          WHERE NAME LIKE 'stage/innodb/buffer%';
   ```

2. Enable the stage event consumer tables, which include `events_stages_current`, `events_stages_history`, and `events_stages_history_long`.

   ```
   mysql> UPDATE performance_schema.setup_consumers SET ENABLED = 'YES'
          WHERE NAME LIKE '%stages%';
   ```

3. Dump the current buffer pool state by enabling `innodb_buffer_pool_dump_now`.

```
mysql> SET GLOBAL innodb_buffer_pool_dump_now=ON;
```

4. Check the buffer pool dump status to ensure that the operation has completed.

```
mysql> SHOW STATUS LIKE 'Innodb_buffer_pool_dump_status'\G
*************************** 1. row ***************************
Variable_name: Innodb_buffer_pool_dump_status
        Value: Buffer pool(s) dump completed at 150202 16:38:58
```

5. Load the buffer pool by enabling `innodb_buffer_pool_load_now`:

```
mysql> SET GLOBAL innodb_buffer_pool_load_now=ON;
```

6. Check the current status of the buffer pool load operation by querying the Performance Schema `events_stages_current` table. The `WORK_COMPLETED` column shows the number of buffer pool pages loaded. The `WORK_ESTIMATED` column provides an estimate of the remaining work, in pages.

```
mysql> SELECT EVENT_NAME, WORK_COMPLETED, WORK_ESTIMATED
       FROM performance_schema.events_stages_current;
+-------------------------------+----------------+----------------+
| EVENT_NAME                    | WORK_COMPLETED | WORK_ESTIMATED |
+-------------------------------+----------------+----------------+
| stage/innodb/buffer pool load |           5353 |           7167 |
+-------------------------------+----------------+----------------+
```

The `events_stages_current` table returns an empty set if the buffer pool load operation has completed. In this case, you can check the `events_stages_history` table to view data for the completed event. For example:

```
mysql> SELECT EVENT_NAME, WORK_COMPLETED, WORK_ESTIMATED
       FROM performance_schema.events_stages_history;
+-------------------------------+----------------+----------------+
| EVENT_NAME                    | WORK_COMPLETED | WORK_ESTIMATED |
+-------------------------------+----------------+----------------+
| stage/innodb/buffer pool load |           7167 |           7167 |
+-------------------------------+----------------+----------------+
```

Note

You can also monitor buffer pool load progress using Performance Schema when loading the buffer pool at startup using `innodb_buffer_pool_load_at_startup`. In this case, the `stage/innodb/buffer pool load` instrument and related consumers must be enabled at startup. For more information, see Section 25.3, “Performance Schema Startup Configuration”.

15.6.3.9 Monitoring the Buffer Pool Using the InnoDB Standard Monitor

`InnoDB` Standard Monitor output, which can be accessed using `SHOW ENGINE INNODB STATUS`, provides metrics that pertain to operation of the `InnoDB` buffer pool. Buffer pool metrics are located in the `BUFFER POOL AND MEMORY` section of `InnoDB` Standard Monitor output and appear similar to the following:

```
----------------------
BUFFER POOL AND MEMORY
----------------------
Total large memory allocated 2198863872
Dictionary memory allocated 776332
Buffer pool size   131072
```

```
Free buffers       124908
Database pages     5720
Old database pages 2071
Modified db pages  910
Pending reads 0
Pending writes: LRU 0, flush list 0, single page 0
Pages made young 4, not young 0
0.10 youngs/s, 0.00 non-youngs/s
Pages read 197, created 5523, written 5060
0.00 reads/s, 190.89 creates/s, 244.94 writes/s
Buffer pool hit rate 1000 / 1000, young-making rate 0 / 1000 not
0 / 1000
Pages read ahead 0.00/s, evicted without access 0.00/s, Random read
ahead 0.00/s
LRU len: 5720, unzip_LRU len: 0
I/O sum[0]:cur[0], unzip sum[0]:cur[0]
```

The following table describes `InnoDB` buffer pool metrics reported by the `InnoDB` Standard Monitor.

Note

Per second averages provided in `InnoDB` Standard Monitor output are based on the elapsed time since `InnoDB` Standard Monitor output was last printed.

Table 15.2 InnoDB Buffer Pool Metrics

Name	Description
Total memory allocated	The total memory allocated for the buffer pool in bytes.
Dictionary memory allocated	The total memory allocated for the `InnoDB` data dictionary in bytes.
Buffer pool size	The total size in pages allocated to the buffer pool.
Free buffers	The total size in pages of the buffer pool free list.
Database pages	The total size in pages of the buffer pool LRU list.
Old database pages	The total size in pages of the buffer pool old LRU sublist.
Modified db pages	The current number of pages modified in the buffer pool.
Pending reads	The number of buffer pool pages waiting to be read in to the buffer pool.
Pending writes LRU	The number of old dirty pages within the buffer pool to be written from the bottom of the LRU list.
Pending writes flush list	The number of buffer pool pages to be flushed during checkpointing.
Pending writes single page	The number of pending independent page writes within the buffer pool.
Pages made young	The total number of pages made young in the buffer pool LRU list (moved to the head of sublist of “new” pages).
Pages made not young	The total number of pages not made young in the buffer pool LRU list (pages that have remained in the “old” sublist without being made young).
youngs/s	The per second average of accesses to old pages in the buffer pool LRU list that have resulted in making pages young. See the notes that follow this table for more information.
non-youngs/s	The per second average of accesses to old pages in the buffer pool LRU list that have resulted in not making pages young. See the notes that follow this table for more information.
Pages read	The total number of pages read from the buffer pool.
Pages created	The total number of pages created within the buffer pool.

Name	Description
Pages written	The total number of pages written from the buffer pool.
reads/s	The per second average number of buffer pool page reads per second.
creates/s	The per second average number of buffer pool pages created per second.
writes/s	The per second average number of buffer pool page writes per second.
Buffer pool hit rate	The buffer pool page hit rate for pages read from the buffer pool memory vs from disk storage.
young-making rate	The average hit rate at which page accesses have resulted in making pages young. See the notes that follow this table for more information.
not (young-making rate)	The average hit rate at which page accesses have not resulted in making pages young. See the notes that follow this table for more information.
Pages read ahead	The per second average of read ahead operations.
Pages evicted without access	The per second average of the pages evicted without being accessed from the buffer pool.
Random read ahead	The per second average of random read ahead operations.
LRU len	The total size in pages of the buffer pool LRU list.
unzip_LRU len	The total size in pages of the buffer pool unzip_LRU list.
I/O sum	The total number of buffer pool LRU list pages accessed, for the last 50 seconds.
I/O cur	The total number of buffer pool LRU list pages accessed.
I/O unzip sum	The total number of buffer pool unzip_LRU list pages accessed.
I/O unzip cur	The total number of buffer pool unzip_LRU list pages accessed.

Notes:

- The `youngs/s` metric only relates to old pages. It is based on the number of accesses to pages and not the number of pages. There can be multiple accesses to a given page, all of which are counted. If you see very low `youngs/s` values when there are no large scans occurring, you might need to reduce the delay time or increase the percentage of the buffer pool used for the old sublist. Increasing the percentage makes the old sublist larger, so pages in that sublist take longer to move to the tail and to be evicted. This increases the likelihood that the pages will be accessed again and be made young.
- The `non-youngs/s` metric only relates to old pages. It is based on the number of accesses to pages and not the number of pages. There can be multiple accesses to a given page, all of which are counted. If you do not see a lot of `non-youngs/s` when you are doing large table scans (and lots of `youngs/s`), increase the delay value.
- The `young-making` rate accounts for accesses to all buffer pool pages, not just accesses to pages in the old sublist. The `young-making` rate and `not` rate do not normally add up to the overall buffer pool hit rate. Page hits in the old sublist cause pages to move to the new sublist, but page hits in the new sublist cause pages to move to the head of the list only if they are a certain distance from the head.
- `not (young-making rate)` is the average hit rate at which page accesses have not resulted in making pages young due to the delay defined by `innodb_old_blocks_time` not being met, or due to page hits in the new sublist that did not result in pages being moved to the head. This rate accounts for accesses to all buffer pool pages, not just accesses to pages in the old sublist.

InnoDB buffer pool server status variables and the INNODB_BUFFER_POOL_STATS table provide many of the same buffer pool metrics found in InnoDB Standard Monitor output. For more information about the INNODB_BUFFER_POOL_STATS table, see Example 15.10, “Querying the INNODB_BUFFER_POOL_STATS Table”.

15.6.4 Configuring InnoDB Change Buffering

When INSERT, UPDATE, and DELETE operations are performed on a table, the values of indexed columns (particularly the values of secondary keys) are often in an unsorted order, requiring substantial I/O to bring secondary indexes up to date. InnoDB has a change buffer that caches changes to secondary index entries when the relevant page is not in the buffer pool, thus avoiding expensive I/O operations by not immediately reading in the page from disk. The buffered changes are merged when the page is loaded to the buffer pool, and the updated page is later flushed to disk. The InnoDB main thread merges buffered changes when the server is nearly idle, and during a slow shutdown.

Because it can result in fewer disk reads and writes, the change buffer feature is most valuable for workloads that are I/O-bound, for example applications with a high volume of DML operations such as bulk inserts.

However, the change buffer occupies a part of the buffer pool, reducing the memory available to cache data pages. If the working set almost fits in the buffer pool, or if your tables have relatively few secondary indexes, it may be useful to disable change buffering. If the working set fits entirely within the buffer, change buffering does not impose extra overhead, because it only applies to pages that are not in the buffer pool.

You can control the extent to which InnoDB performs change buffering using the innodb_change_buffering configuration parameter. You can enable or disable buffering for inserts, delete operations (when index records are initially marked for deletion) and purge operations (when index records are physically deleted). An update operation is a combination of an insert and a delete. The default innodb_change_buffering value is all.

Permitted innodb_change_buffering values include:

- **all**

  The default value: buffer inserts, delete-marking operations, and purges.

- **none**

  Do not buffer any operations.

- **inserts**

  Buffer insert operations.

- **deletes**

  Buffer delete-marking operations.

- **changes**

  Buffer both inserts and delete-marking operations.

- **purges**

  Buffer the physical deletion operations that happen in the background.

You can set the innodb_change_buffering parameter in the MySQL option file (my.cnf or my.ini) or change it dynamically with the SET GLOBAL command, which requires the SYSTEM_VARIABLES_ADMIN or SUPER privilege. Changing the setting affects the buffering of new operations; the merging of existing buffered entries is not affected.

Change buffering is not supported for a secondary index if the index contains a descending index column or if the primary key includes a descending index column.

For related information, see Section 15.4.2, “Change Buffer”. For information about configuring change buffer size, see Section 15.6.4.1, “Configuring the Change Buffer Maximum Size”.

15.6.4.1 Configuring the Change Buffer Maximum Size

The `innodb_change_buffer_max_size` configuration option allows you to configure the maximum size of the change buffer as a percentage of the total size of the buffer pool. By default, `innodb_change_buffer_max_size` is set to 25. The maximum setting is 50.

You might consider increasing `innodb_change_buffer_max_size` on a MySQL server with heavy insert, update, and delete activity, where change buffer merging does not keep pace with new change buffer entries, causing the change buffer to reach its maximum size limit.

You might consider decreasing `innodb_change_buffer_max_size` on a MySQL server with static data used for reporting, or if the change buffer consumes too much of the memory space that is shared with the buffer pool, causing pages to age out of the buffer pool sooner than desired.

Test different settings with a representative workload to determine an optimal configuration. The `innodb_change_buffer_max_size` setting is dynamic, which allows you modify the setting without restarting the server.

15.6.5 Configuring Thread Concurrency for InnoDB

`InnoDB` uses operating system threads to process requests from user transactions. (Transactions may issue many requests to `InnoDB` before they commit or roll back.) On modern operating systems and servers with multi-core processors, where context switching is efficient, most workloads run well without any limit on the number of concurrent threads.

In situations where it is helpful to minimize context switching between threads, `InnoDB` can use a number of techniques to limit the number of concurrently executing operating system threads (and thus the number of requests that are processed at any one time). When `InnoDB` receives a new request from a user session, if the number of threads concurrently executing is at a pre-defined limit, the new request sleeps for a short time before it tries again. A request that cannot be rescheduled after the sleep is put in a first-in/first-out queue and eventually is processed. Threads waiting for locks are not counted in the number of concurrently executing threads.

You can limit the number of concurrent threads by setting the configuration parameter `innodb_thread_concurrency`. Once the number of executing threads reaches this limit, additional threads sleep for a number of microseconds, set by the configuration parameter `innodb_thread_sleep_delay`, before being placed into the queue.

You can set the configuration option `innodb_adaptive_max_sleep_delay` to the highest value you would allow for `innodb_thread_sleep_delay`, and `InnoDB` automatically adjusts `innodb_thread_sleep_delay` up or down depending on the current thread-scheduling activity. This dynamic adjustment helps the thread scheduling mechanism to work smoothly during times when the system is lightly loaded and when it is operating near full capacity.

The default value for `innodb_thread_concurrency` and the implied default limit on the number of concurrent threads has been changed in various releases of MySQL and `InnoDB`. The default value of `innodb_thread_concurrency` is `0`, so that by default there is no limit on the number of concurrently executing threads.

`InnoDB` causes threads to sleep only when the number of concurrent threads is limited. When there is no limit on the number of threads, all contend equally to be scheduled. That is, if `innodb_thread_concurrency` is `0`, the value of `innodb_thread_sleep_delay` is ignored.

When there is a limit on the number of threads (when `innodb_thread_concurrency` is > 0), `InnoDB` reduces context switching overhead by permitting multiple requests made during

the execution of a *single SQL statement* to enter InnoDB without observing the limit set by innodb_thread_concurrency. Since an SQL statement (such as a join) may comprise multiple row operations within InnoDB, InnoDB assigns a specified number of "tickets" that allow a thread to be scheduled repeatedly with minimal overhead.

When a new SQL statement starts, a thread has no tickets, and it must observe innodb_thread_concurrency. Once the thread is entitled to enter InnoDB, it is assigned a number of tickets that it can use for subsequently entering InnoDB to perform row operations. If the tickets run out, the thread is evicted, and innodb_thread_concurrency is observed again which may place the thread back into the first-in/first-out queue of waiting threads. When the thread is once again entitled to enter InnoDB, tickets are assigned again. The number of tickets assigned is specified by the global option innodb_concurrency_tickets, which is 5000 by default. A thread that is waiting for a lock is given one ticket once the lock becomes available.

The correct values of these variables depend on your environment and workload. Try a range of different values to determine what value works for your applications. Before limiting the number of concurrently executing threads, review configuration options that may improve the performance of InnoDB on multi-core and multi-processor computers, such as innodb_adaptive_hash_index.

For general performance information about MySQL thread handling, see Section 8.12.4.1, "How MySQL Uses Threads for Client Connections".

15.6.6 Configuring the Number of Background InnoDB I/O Threads

InnoDB uses background threads to service various types of I/O requests. You can configure the number of background threads that service read and write I/O on data pages using the innodb_read_io_threads and innodb_write_io_threads configuration parameters. These parameters signify the number of background threads used for read and write requests, respectively. They are effective on all supported platforms. You can set values for these parameters in the MySQL option file (my.cnf or my.ini); you cannot change values dynamically. The default value for these parameters is 4 and permissible values range from 1-64.

The purpose of these configuration options to make InnoDB more scalable on high end systems. Each background thread can handle up to 256 pending I/O requests. A major source of background I/O is read-ahead requests. InnoDB tries to balance the load of incoming requests in such way that most background threads share work equally. InnoDB also attempts to allocate read requests from the same extent to the same thread, to increase the chances of coalescing the requests. If you have a high end I/O subsystem and you see more than 64 × innodb_read_io_threads pending read requests in SHOW ENGINE INNODB STATUS output, you might improve performance by increasing the value of innodb_read_io_threads.

On Linux systems, InnoDB uses the asynchronous I/O subsystem by default to perform read-ahead and write requests for data file pages, which changes the way that InnoDB background threads service these types of I/O requests. For more information, see Section 15.6.7, "Using Asynchronous I/O on Linux".

For more information about InnoDB I/O performance, see Section 8.5.8, "Optimizing InnoDB Disk I/O".

15.6.7 Using Asynchronous I/O on Linux

InnoDB uses the asynchronous I/O subsystem (native AIO) on Linux to perform readahead and write requests for data file pages. This behavior is controlled by the innodb_use_native_aio configuration option, which applies to Linux systems only and is enabled by default. On other Unix-like systems, InnoDB uses synchronous I/O only. Historically, InnoDB only used asynchronous I/O on Windows systems. Using the asynchronous I/O subsystem on Linux requires the libaio library.

With synchronous I/O, query threads queue I/O requests, and InnoDB background threads retrieve the queued requests one at a time, issuing a synchronous I/O call for each. When an I/O request is completed and the I/O call returns, the InnoDB background thread that is handling the request

calls an I/O completion routine and returns to process the next request. The number of requests that can be processed in parallel is *n*, where *n* is the number of `InnoDB` background threads. The number of `InnoDB` background threads is controlled by `innodb_read_io_threads` and `innodb_write_io_threads`. See Section 15.6.6, "Configuring the Number of Background InnoDB I/O Threads".

With native AIO, query threads dispatch I/O requests directly to the operating system, thereby removing the limit imposed by the number of background threads. `InnoDB` background threads wait for I/O events to signal completed requests. When a request is completed, a background thread calls an I/O completion routine and resumes waiting for I/O events.

The advantage of native AIO is scalability for heavily I/O-bound systems that typically show many pending reads/writes in `SHOW ENGINE INNODB STATUS\G` output. The increase in parallel processing when using native AIO means that the type of I/O scheduler or properties of the disk array controller have a greater influence on I/O performance.

A potential disadvantage of native AIO for heavily I/O-bound systems is lack of control over the number of I/O write requests dispatched to the operating system at once. Too many I/O write requests dispatched to the operating system for parallel processing could, in some cases, result in I/O read starvation, depending on the amount of I/O activity and system capabilities.

If a problem with the asynchronous I/O subsystem in the OS prevents `InnoDB` from starting, you can start the server with `innodb_use_native_aio=0`. This option may also be disabled automatically during startup if `InnoDB` detects a potential problem such as a combination of `tmpdir` location, `tmpfs` file system, and Linux kernel that does not support asynchronous I/O on `tmpfs`.

15.6.8 Configuring the InnoDB Master Thread I/O Rate

The master thread in InnoDB is a thread that performs various tasks in the background. Most of these tasks are I/O related, such as flushing dirty pages from the buffer pool or writing changes from the insert buffer to the appropriate secondary indexes. The master thread attempts to perform these tasks in a way that does not adversely affect the normal working of the server. It tries to estimate the free I/O bandwidth available and tune its activities to take advantage of this free capacity. Historically, InnoDB has used a hard coded value of 100 IOPs (input/output operations per second) as the total I/O capacity of the server.

The parameter `innodb_io_capacity` indicates the overall I/O capacity available to InnoDB. This parameter should be set to approximately the number of I/O operations that the system can perform per second. The value depends on your system configuration. When `innodb_io_capacity` is set, the master threads estimates the I/O bandwidth available for background tasks based on the set value. Setting the value to `100` reverts to the old behavior.

You can set the value of `innodb_io_capacity` to any number 100 or greater. The default value is `200`, reflecting that the performance of typical modern I/O devices is higher than in the early days of MySQL. Typically, values around the previous default of 100 are appropriate for consumer-level storage devices, such as hard drives up to 7200 RPMs. Faster hard drives, RAID configurations, and SSDs benefit from higher values.

The `innodb_io_capacity` setting is a total limit for all buffer pool instances. When dirty pages are flushed, the `innodb_io_capacity` limit is divided equally among buffer pool instances. For more information, see the `innodb_io_capacity` system variable description.

You can set the value of this parameter in the MySQL option file (`my.cnf` or `my.ini`) or change it dynamically with the `SET GLOBAL` command, which requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege.

The `innodb_flush_sync` configuration option causes the `innodb_io_capacity` setting to be ignored during bursts of I/O activity that occur at checkpoints. `innodb_flush_sync` is enabled by default.

In earlier MySQL releases, the InnoDB master thread also performed any needed purge operations. Those I/O operations are now performed by other background threads, whose number is controlled by the innodb_purge_threads configuration option.

For more information about InnoDB I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

15.6.9 Configuring Spin Lock Polling

Many InnoDB mutexes and rw-locks are reserved for a short time. On a multi-core system, it can be more efficient for a thread to continuously check if it can acquire a mutex or rw-lock for a while before sleeping. If the mutex or rw-lock becomes available during this polling period, the thread can continue immediately, in the same time slice. However, too-frequent polling by multiple threads of a shared object can cause “cache ping pong”, different processors invalidating portions of each others' cache. InnoDB minimizes this issue by waiting a random time between subsequent polls. The delay is implemented as a busy loop.

You can control the maximum delay between testing a mutex or rw-lock using the parameter innodb_spin_wait_delay. The duration of the delay loop depends on the C compiler and the target processor. (In the 100MHz Pentium era, the unit of delay was one microsecond.) On a system where all processor cores share a fast cache memory, you might reduce the maximum delay or disable the busy loop altogether by setting innodb_spin_wait_delay=0. On a system with multiple processor chips, the effect of cache invalidation can be more significant and you might increase the maximum delay.

The default value of innodb_spin_wait_delay is 6. The spin wait delay is a dynamic, global parameter that you can specify in the MySQL option file (my.cnf or my.ini) or change at runtime with the command SET GLOBAL innodb_spin_wait_delay=*delay*, where *delay* is the desired maximum delay. Changing the setting requires the SYSTEM_VARIABLES_ADMIN or SUPER privilege.

For performance considerations for InnoDB locking operations, see Section 8.11, “Optimizing Locking Operations”.

15.6.10 Configuring InnoDB Purge Scheduling

The purge operations (a type of garbage collection) that InnoDB performs automatically may be performed by one or more separate threads rather than as part of the master thread. The use of separate threads improves scalability by allowing the main database operations to run independently from maintenance work happening in the background.

To control this feature, increase the value of the configuration option innodb_purge_threads. If DML action is concentrated on a single table or a few tables, keep the setting low so that the threads do not contend with each other for access to the busy tables. If DML operations are spread across many tables, increase the setting. Its maximum is 32. innodb_purge_threads is a non-dynamic configuration option, which means it cannot be configured at runtime.

There is another related configuration option, innodb_purge_batch_size with a default value of 300 and maximum value of 5000. This option is mainly intended for experimentation and tuning of purge operations, and should not be interesting to typical users.

For more information about InnoDB I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

15.6.11 Configuring Optimizer Statistics for InnoDB

This section describes how to configure persistent and non-persistent optimizer statistics for InnoDB tables.

Persistent optimizer statistics are persisted across server restarts, allowing for greater plan stability and more consistent query performance. Persistent optimizer statistics also provide control and flexibility with these additional benefits:

- You can use the `innodb_stats_auto_recalc` configuration option to control whether statistics are updated automatically after substantial changes to a table.
- You can use the `STATS_PERSISTENT`, `STATS_AUTO_RECALC`, and `STATS_SAMPLE_PAGES` clauses with `CREATE TABLE` and `ALTER TABLE` statements to configure optimizer statistics for individual tables.
- You can query optimizer statistics data in the `mysql.innodb_table_stats` and `mysql.innodb_index_stats` tables.
- You can view the `last_update` column of the `mysql.innodb_table_stats` and `mysql.innodb_index_stats` tables to see when statistics were last updated.
- You can manually modify the `mysql.innodb_table_stats` and `mysql.innodb_index_stats` tables to force a specific query optimization plan or to test alternative plans without modifying the database.

The persistent optimizer statistics feature is enabled by default (`innodb_stats_persistent=ON`).

Non-persistent optimizer statistics are cleared on each server restart and after some other operations, and recomputed on the next table access. As a result, different estimates could be produced when recomputing statistics, leading to different choices in execution plans and variations in query performance.

This section also provides information about estimating `ANALYZE TABLE` complexity, which may be useful when attempting to achieve a balance between accurate statistics and `ANALYZE TABLE` execution time.

15.6.11.1 Configuring Persistent Optimizer Statistics Parameters

The persistent optimizer statistics feature improves plan stability by storing statistics to disk and making them persistent across server restarts so that the optimizer is more likely to make consistent choices each time for a given query.

Optimizer statistics are persisted to disk when `innodb_stats_persistent=ON` or when individual tables are created or altered with `STATS_PERSISTENT=1`. `innodb_stats_persistent` is enabled by default.

Formerly, optimizer statistics were cleared on each server restart and after some other operations, and recomputed on the next table access. Consequently, different estimates could be produced when recalculating statistics, leading to different choices in query execution plans and thus variations in query performance.

Persistent statistics are stored in the `mysql.innodb_table_stats` and `mysql.innodb_index_stats` tables, as described in InnoDB Persistent Statistics Tables.

To revert to using non-persistent optimizer statistics, you can modify tables using an `ALTER TABLE tbl_name STATS_PERSISTENT=0` statement. For related information, see Section 15.6.11.2, “Configuring Non-Persistent Optimizer Statistics Parameters”

Configuring Automatic Statistics Calculation for Persistent Optimizer Statistics

The `innodb_stats_auto_recalc` configuration option, which is enabled by default, determines whether statistics are calculated automatically whenever a table undergoes substantial changes (to more than 10% of the rows). You can also configure automatic statistics recalculation for individual tables using a `STATS_AUTO_RECALC` clause in a `CREATE TABLE` or `ALTER TABLE` statement. `innodb_stats_auto_recalc` is enabled by default.

Because of the asynchronous nature of automatic statistics recalculation (which occurs in the background), statistics may not be recalculated instantly after running a DML operation that affects more than 10% of a table, even when `innodb_stats_auto_recalc` is enabled. In

some cases, statistics recalculation may be delayed by a few seconds. If up-to-date statistics are required immediately after changing significant portions of a table, run ANALYZE TABLE to initiate a synchronous (foreground) recalculation of statistics.

If innodb_stats_auto_recalc is disabled, ensure the accuracy of optimizer statistics by issuing the ANALYZE TABLE statement for each applicable table after making substantial changes to indexed columns. You might run this statement in your setup scripts after representative data has been loaded into the table, and run it periodically after DML operations significantly change the contents of indexed columns, or on a schedule at times of low activity. When a new index is added to an existing table, index statistics are calculated and added to the innodb_index_stats table regardless of the value of innodb_stats_auto_recalc.

Caution

To ensure statistics are gathered when a new index is created, either enable the innodb_stats_auto_recalc option, or run ANALYZE TABLE after creating each new index when the persistent statistics mode is enabled.

Configuring Optimizer Statistics Parameters for Individual Tables

innodb_stats_persistent, innodb_stats_auto_recalc, and innodb_stats_persistent_sample_pages are global configuration options. To override these system-wide settings and configure optimizer statistics parameters for individual tables, you can define STATS_PERSISTENT, STATS_AUTO_RECALC, and STATS_SAMPLE_PAGES clauses in CREATE TABLE or ALTER TABLE statements.

- STATS_PERSISTENT specifies whether to enable persistent statistics for an InnoDB table. The value DEFAULT causes the persistent statistics setting for the table to be determined by the innodb_stats_persistent configuration option. The value 1 enables persistent statistics for the table, while the value 0 turns off this feature. After enabling persistent statistics through a CREATE TABLE or ALTER TABLE statement, issue an ANALYZE TABLE statement to calculate the statistics, after loading representative data into the table.
- STATS_AUTO_RECALC specifies whether to automatically recalculate persistent statistics for an InnoDB table. The value DEFAULT causes the persistent statistics setting for the table to be determined by the innodb_stats_auto_recalc configuration option. The value 1 causes statistics to be recalculated when 10% of the data in the table has changed. The value 0 prevents automatic recalculation for this table; with this setting, issue an ANALYZE TABLE statement to recalculate the statistics after making substantial changes to the table.
- STATS_SAMPLE_PAGES specifies the number of index pages to sample when estimating cardinality and other statistics for an indexed column, such as those calculated by ANALYZE TABLE.

All three clauses are specified in the following CREATE TABLE example:

```
CREATE TABLE `t1` (
`id` int(8) NOT NULL auto_increment,
`data` varchar(255),
`date` datetime,
PRIMARY KEY  (`id`),
INDEX `DATE_IX` (`date`)
) ENGINE=InnoDB,
  STATS_PERSISTENT=1,
  STATS_AUTO_RECALC=1,
  STATS_SAMPLE_PAGES=25;
```

Configuring the Number of Sampled Pages for InnoDB Optimizer Statistics

The MySQL query optimizer uses estimated statistics about key distributions to choose the indexes for an execution plan, based on the relative selectivity of the index. Operations such as ANALYZE TABLE cause InnoDB to sample random pages from each index on a table to estimate the cardinality of the index. (This technique is known as random dives.)

To give you control over the quality of the statistics estimate (and thus better information for the query optimizer), you can change the number of sampled pages using the parameter `innodb_stats_persistent_sample_pages`, which can be set at runtime.

`innodb_stats_persistent_sample_pages` has a default value of 20. As a general guideline, consider modifying this parameter when encountering the following issues:

1. *Statistics are not accurate enough and the optimizer chooses suboptimal plans*, as shown by `EXPLAIN` output. The accuracy of statistics can be checked by comparing the actual cardinality of an index (as returned by running `SELECT DISTINCT` on the index columns) with the estimates provided in the `mysql.innodb_index_stats` persistent statistics table.

   If it is determined that statistics are not accurate enough, the value of `innodb_stats_persistent_sample_pages` should be increased until the statistics estimates are sufficiently accurate. Increasing `innodb_stats_persistent_sample_pages` too much, however, could cause `ANALYZE TABLE` to run slowly.

2. *`ANALYZE TABLE` is too slow*. In this case `innodb_stats_persistent_sample_pages` should be decreased until `ANALYZE TABLE` execution time is acceptable. Decreasing the value too much, however, could lead to the first problem of inaccurate statistics and suboptimal query execution plans.

   If a balance cannot be achieved between accurate statistics and `ANALYZE TABLE` execution time, consider decreasing the number of indexed columns in the table or limiting the number of partitions to reduce `ANALYZE TABLE` complexity. The number of columns in the table's primary key is also important to consider, as primary key columns are appended to each non-unique index.

   For related information, see Section 15.6.11.3, “Estimating ANALYZE TABLE Complexity for InnoDB Tables”.

Including Delete-marked Records in Persistent Statistics Calculations

By default, `InnoDB` reads uncommitted data when calculating statistics. In the case of an uncommitted transaction that deletes rows from a table, `InnoDB` excludes records that are delete-marked when calculating row estimates and index statistics, which can lead to non-optimal execution plans for other transactions that are operating on the table concurrently using a transaction isolation level other than `READ UNCOMMITTED`. To avoid this scenario, `innodb_stats_include_delete_marked` can be enabled to ensure that `InnoDB` includes delete-marked records when calculating persistent optimizer statistics.

When `innodb_stats_include_delete_marked` is enabled, `ANALYZE TABLE` considers delete-marked records when recalculating statistics.

`innodb_stats_include_delete_marked` is a global setting that affects all `InnoDB` tables, and it is only applicable to persistent optimizer statistics.

InnoDB Persistent Statistics Tables

The persistent statistics feature relies on the internally managed tables in the `mysql` database, named `innodb_table_stats` and `innodb_index_stats`. These tables are set up automatically in all install, upgrade, and build-from-source procedures.

Table 15.3 Columns of innodb_table_stats

Column name	Description
`database_name`	Database name
`table_name`	Table name, partition name, or subpartition name
`last_update`	A timestamp indicating the last time that `InnoDB` updated this row

Column name	Description
n_rows	The number of rows in the table
clustered_index_size	The size of the primary index, in pages
sum_of_other_index_sizes	The total size of other (non-primary) indexes, in pages

Table 15.4 Columns of innodb_index_stats

Column name	Description
database_name	Database name
table_name	Table name, partition name, or subpartition name
index_name	Index name
last_update	A timestamp indicating the last time that InnoDB updated this row
stat_name	The name of the statistic, whose value is reported in the stat_value column
stat_value	The value of the statistic that is named in stat_name column
sample_size	The number of pages sampled for the estimate provided in the stat_value column
stat_description	Description of the statistic that is named in the stat_name column

Both the innodb_table_stats and innodb_index_stats tables include a last_update column showing when InnoDB last updated index statistics, as shown in the following example:

```
mysql> select * from innodb_table_stats \G
*************************** 1. row ***************************
           database_name: sakila
              table_name: actor
             last_update: 2014-05-28 16:16:44
                 n_rows: 200
    clustered_index_size: 1
sum_of_other_index_sizes: 1
...
```

```
mysql> select * from innodb_index_stats \G
*************************** 1. row ***************************
   database_name: sakila
      table_name: actor
      index_name: PRIMARY
     last_update: 2014-05-28 16:16:44
       stat_name: n_diff_pfx01
      stat_value: 200
     sample_size: 1
     ...
```

The innodb_table_stats and innodb_index_stats tables are ordinary tables and can be updated manually. The ability to update statistics manually makes it possible to force a specific query optimization plan or test alternative plans without modifying the database. If you manually update statistics, issue the FLUSH TABLE *tbl_name* command to make MySQL reload the updated statistics.

InnoDB Persistent Statistics Tables Example

The innodb_table_stats table contains one row per table. The data collected is demonstrated in the following example.

Table t1 contains a primary index (columns a, b) secondary index (columns c, d), and unique index (columns e, f):

```
CREATE TABLE t1 (
a INT, b INT, c INT, d INT, e INT, f INT,
```

```
PRIMARY KEY (a, b), KEY i1 (c, d), UNIQUE KEY i2uniq (e, f)
) ENGINE=INNODB;
```

After inserting five rows of sample data, the table appears as follows:

```
mysql> SELECT * FROM t1;
+---+---+------+------+------+------+
| a | b | c    | d    | e    | f    |
+---+---+------+------+------+------+
| 1 | 1 |   10 |   11 |  100 |  101 |
| 1 | 2 |   10 |   11 |  200 |  102 |
| 1 | 3 |   10 |   11 |  100 |  103 |
| 1 | 4 |   10 |   12 |  200 |  104 |
| 1 | 5 |   10 |   12 |  100 |  105 |
+---+---+------+------+------+------+
```

To immediately update statistics, run `ANALYZE TABLE` (if `innodb_stats_auto_recalc` is enabled, statistics are updated automatically within a few seconds assuming that the 10% threshold for changed table rows is reached):

```
mysql> ANALYZE TABLE t1;
+---------+---------+----------+----------+
| Table   | Op      | Msg_type | Msg_text |
+---------+---------+----------+----------+
| test.t1 | analyze | status   | OK       |
+---------+---------+----------+----------+
```

Table statistics for table `t1` show the last time `InnoDB` updated the table statistics (`2014-03-14 14:36:34`), the number of rows in the table (`5`), the clustered index size (`1` page), and the combined size of the other indexes (`2` pages).

```
mysql> SELECT * FROM mysql.innodb_table_stats WHERE table_name like 't1'\G
*************************** 1. row ***************************
           database_name: test
              table_name: t1
             last_update: 2014-03-14 14:36:34
                 n_rows: 5
    clustered_index_size: 1
sum_of_other_index_sizes: 2
```

The `innodb_index_stats` table contains multiple rows for each index. Each row in the `innodb_index_stats` table provides data related to a particular index statistic which is named in the `stat_name` column and described in the `stat_description` column. For example:

```
mysql> SELECT index_name, stat_name, stat_value, stat_description
    -> FROM mysql.innodb_index_stats WHERE table_name like 't1';
+------------+--------------+------------+----------------------------------+
| index_name | stat_name    | stat_value | stat_description                 |
+------------+--------------+------------+----------------------------------+
| PRIMARY    | n_diff_pfx01 |          1 | a                                |
| PRIMARY    | n_diff_pfx02 |          5 | a,b                              |
| PRIMARY    | n_leaf_pages |          1 | Number of leaf pages in the index |
| PRIMARY    | size         |          1 | Number of pages in the index      |
| i1         | n_diff_pfx01 |          1 | c                                |
| i1         | n_diff_pfx02 |          2 | c,d                              |
| i1         | n_diff_pfx03 |          2 | c,d,a                            |
| i1         | n_diff_pfx04 |          5 | c,d,a,b                          |
| i1         | n_leaf_pages |          1 | Number of leaf pages in the index |
| i1         | size         |          1 | Number of pages in the index      |
| i2uniq     | n_diff_pfx01 |          2 | e                                |
| i2uniq     | n_diff_pfx02 |          5 | e,f                              |
| i2uniq     | n_leaf_pages |          1 | Number of leaf pages in the index |
| i2uniq     | size         |          1 | Number of pages in the index      |
+------------+--------------+------------+----------------------------------+
```

The `stat_name` column shows the following types of statistics:

- `size`: Where `stat_name=size`, the `stat_value` column displays the total number of pages in the index.

- `n_leaf_pages`: Where `stat_name=n_leaf_pages`, the `stat_value` column displays the number of leaf pages in the index.
- `n_diff_pfxNN`: Where `stat_name=n_diff_pfx01`, the `stat_value` column displays the number of distinct values in the first column of the index. Where `stat_name=n_diff_pfx02`, the `stat_value` column displays the number of distinct values in the first two columns of the index, and so on. Additionally, where `stat_name=n_diff_pfxNN`, the `stat_description` column shows a comma separated list of the index columns that are counted.

To further illustrate the `n_diff_pfxNN` statistic, which provides cardinality data, consider the `t1` table example. As shown below, the `t1` table is created with a primary index (columns `a`, `b`), a secondary index (columns `c`, `d`), and a unique index (columns `e`, `f`):

```
CREATE TABLE t1 (
  a INT, b INT, c INT, d INT, e INT, f INT,
  PRIMARY KEY (a, b), KEY i1 (c, d), UNIQUE KEY i2uniq (e, f)
) ENGINE=INNODB;
```

After inserting five rows of sample data, the table appears as follows:

```
mysql> SELECT * FROM t1;
+---+---+------+------+------+------+
| a | b | c    | d    | e    | f    |
+---+---+------+------+------+------+
| 1 | 1 |   10 |   11 |  100 |  101 |
| 1 | 2 |   10 |   11 |  200 |  102 |
| 1 | 3 |   10 |   11 |  100 |  103 |
| 1 | 4 |   10 |   12 |  200 |  104 |
| 1 | 5 |   10 |   12 |  100 |  105 |
+---+---+------+------+------+------+
```

When you query the `index_name`, `stat_name`, `stat_value`, and `stat_description` where `stat_name LIKE 'n_diff%'`, the following result set is returned:

```
mysql> SELECT index_name, stat_name, stat_value, stat_description
    -> FROM mysql.innodb_index_stats
    -> WHERE table_name like 't1' AND stat_name LIKE 'n_diff%';
+------------+--------------+------------+------------------+
| index_name | stat_name    | stat_value | stat_description |
+------------+--------------+------------+------------------+
| PRIMARY    | n_diff_pfx01 |          1 | a                |
| PRIMARY    | n_diff_pfx02 |          5 | a,b              |
| i1         | n_diff_pfx01 |          1 | c                |
| i1         | n_diff_pfx02 |          2 | c,d              |
| i1         | n_diff_pfx03 |          2 | c,d,a            |
| i1         | n_diff_pfx04 |          5 | c,d,a,b          |
| i2uniq     | n_diff_pfx01 |          2 | e                |
| i2uniq     | n_diff_pfx02 |          5 | e,f              |
+------------+--------------+------------+------------------+
```

For the `PRIMARY` index, there are two `n_diff%` rows. The number of rows is equal to the number of columns in the index.

Note

For non-unique indexes, `InnoDB` appends the columns of the primary key.

- Where `index_name=PRIMARY` and `stat_name=n_diff_pfx01`, the `stat_value` is `1`, which indicates that there is a single distinct value in the first column of the index (column `a`). The number of distinct values in column `a` is confirmed by viewing the data in column `a` in table `t1`, in which there is a single distinct value (`1`). The counted column (`a`) is shown in the `stat_description` column of the result set.
- Where `index_name=PRIMARY` and `stat_name=n_diff_pfx02`, the `stat_value` is `5`, which indicates that there are five distinct values in the two columns of the index (`a,b`). The number of distinct values in columns `a` and `b` is confirmed by viewing the data in columns `a` and `b` in table `t1`,

in which there are five distinct values: (1,1), (1,2), (1,3), (1,4) and (1,5). The counted columns (a,b) are shown in the stat_description column of the result set.

For the secondary index (i1), there are four n_diff% rows. Only two columns are defined for the secondary index (c,d) but there are four n_diff% rows for the secondary index because InnoDB suffixes all non-unique indexes with the primary key. As a result, there are four n_diff% rows instead of two to account for the both the secondary index columns (c,d) and the primary key columns (a,b).

- Where index_name=i1 and stat_name=n_diff_pfx01, the stat_value is 1, which indicates that there is a single distinct value in the first column of the index (column c). The number of distinct values in column c is confirmed by viewing the data in column c in table t1, in which there is a single distinct value: (10). The counted column (c) is shown in the stat_description column of the result set.
- Where index_name=i1 and stat_name=n_diff_pfx02, the stat_value is 2, which indicates that there are two distinct values in the first two columns of the index (c,d). The number of distinct values in columns c an d is confirmed by viewing the data in columns c and d in table t1, in which there are two distinct values: (10,11) and (10,12). The counted columns (c,d) are shown in the stat_description column of the result set.
- Where index_name=i1 and stat_name=n_diff_pfx03, the stat_value is 2, which indicates that there are two distinct values in the first three columns of the index (c,d,a). The number of distinct values in columns c, d, and a is confirmed by viewing the data in column c, d, and a in table t1, in which there are two distinct values: (10,11,1) and (10,12,1). The counted columns (c,d,a) are shown in the stat_description column of the result set.
- Where index_name=i1 and stat_name=n_diff_pfx04, the stat_value is 5, which indicates that there are five distinct values in the four columns of the index (c,d,a,b). The number of distinct values in columns c, d, a and b is confirmed by viewing the data in columns c, d, a, and b in table t1, in which there are five distinct values: (10,11,1,1), (10,11,1,2), (10,11,1,3), (10,12,1,4) and (10,12,1,5). The counted columns (c,d,a,b) are shown in the stat_description column of the result set.

For the unique index (i2uniq), there are two n_diff% rows.

- Where index_name=i2uniq and stat_name=n_diff_pfx01, the stat_value is 2, which indicates that there are two distinct values in the first column of the index (column e). The number of distinct values in column e is confirmed by viewing the data in column e in table t1, in which there are two distinct values: (100) and (200). The counted column (e) is shown in the stat_description column of the result set.
- Where index_name=i2uniq and stat_name=n_diff_pfx02, the stat_value is 5, which indicates that there are five distinct values in the two columns of the index (e,f). The number of distinct values in columns e and f is confirmed by viewing the data in columns e and f in table t1, in which there are five distinct values: (100,101), (200,102), (100,103), (200,104) and (100,105). The counted columns (e,f) are shown in the stat_description column of the result set.

Retrieving Index Size Using the innodb_index_stats Table

The size of indexes for tables, partitions, or subpartitions can be retrieved using the innodb_index_stats table. In the following example, index sizes are retrieved for table t1. For a definition of table t1 and corresponding index statistics, see InnoDB Persistent Statistics Tables Example.

```
mysql> SELECT SUM(stat_value) pages, index_name,
    -> SUM(stat_value)*@@innodb_page_size size
    -> FROM mysql.innodb_index_stats WHERE table_name='t1'
    -> AND stat_name = 'size' GROUP BY index_name;
+-------+------------+-------+
| pages | index_name | size  |
+-------+------------+-------+
|     1 | PRIMARY    | 16384 |
```

```
|     1 | i1         | 16384 |
|     1 | i2uniq     | 16384 |
+-------+------------+-------+
```

For partitions or subpartitions, the same query with a modified `WHERE` clause can be used to retrieve index sizes. For example, the following query retrieves index sizes for partitions of table `t1`:

```
mysql> SELECT SUM(stat_value) pages, index_name,
    -> SUM(stat_value)*@@innodb_page_size size
    -> FROM mysql.innodb_index_stats WHERE table_name like 't1#P%'
    -> AND stat_name = 'size' GROUP BY index_name;
```

15.6.11.2 Configuring Non-Persistent Optimizer Statistics Parameters

This section describes how to configure non-persistent optimizer statistics. Optimizer statistics are not persisted to disk when `innodb_stats_persistent=OFF` or when individual tables are created or altered with `STATS_PERSISTENT=0`. Instead, statistics are stored in memory, and are lost when the server is shut down. Statistics are also updated periodically by certain operations and under certain conditions.

Optimizer statistics are persisted to disk by default, enabled by the `innodb_stats_persistent` configuration option. For information about persistent optimizer statistics, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

Optimizer Statistics Updates

Non-persistent optimizer statistics are updated when:

- Running `ANALYZE TABLE`.
- Running `SHOW TABLE STATUS`, `SHOW INDEX`, or querying the `INFORMATION_SCHEMA.TABLES` or `INFORMATION_SCHEMA.STATISTICS` tables with the `innodb_stats_on_metadata` option enabled.

  The default setting for `innodb_stats_on_metadata` is `OFF`. Enabling `innodb_stats_on_metadata` may reduce access speed for schemas that have a large number of tables or indexes, and reduce stability of execution plans for queries that involve `InnoDB` tables. `innodb_stats_on_metadata` is configured globally using a `SET` statement.

  ```
  SET GLOBAL innodb_stats_on_metadata=ON
  ```

  Note

  `innodb_stats_on_metadata` only applies when optimizer statistics are configured to be non-persistent (when `innodb_stats_persistent` is disabled).

- Starting a `mysql` client with the `--auto-rehash` option enabled, which is the default. The `auto-rehash` option causes all `InnoDB` tables to be opened, and the open table operations cause statistics to be recalculated.

  To improve the start up time of the `mysql` client and to updating statistics, you can turn off `auto-rehash` using the `--disable-auto-rehash` option. The `auto-rehash` feature enables automatic name completion of database, table, and column names for interactive users.

- A table is first opened.
- `InnoDB` detects that 1 / 16 of table has been modified since the last time statistics were updated.

Configuring the Number of Sampled Pages

The MySQL query optimizer uses estimated statistics about key distributions to choose the indexes for an execution plan, based on the relative selectivity of the index. When `InnoDB` updates optimizer

statistics, it samples random pages from each index on a table to estimate the cardinality of the index. (This technique is known as random dives.)

To give you control over the quality of the statistics estimate (and thus better information for the query optimizer), you can change the number of sampled pages using the parameter `innodb_stats_transient_sample_pages`. The default number of sampled pages is 8, which could be insufficient to produce an accurate estimate, leading to poor index choices by the query optimizer. This technique is especially important for large tables and tables used in joins. Unnecessary full table scans for such tables can be a substantial performance issue. See Section 8.2.1.19, "Avoiding Full Table Scans" for tips on tuning such queries. `innodb_stats_transient_sample_pages` is a global parameter that can be set at runtime.

The value of `innodb_stats_transient_sample_pages` affects the index sampling for all `InnoDB` tables and indexes when `innodb_stats_persistent=0`. Be aware of the following potentially significant impacts when you change the index sample size:

- Small values like 1 or 2 can result in inaccurate estimates of cardinality.
- Increasing the `innodb_stats_transient_sample_pages` value might require more disk reads. Values much larger than 8 (say, 100), can cause a significant slowdown in the time it takes to open a table or execute `SHOW TABLE STATUS`.
- The optimizer might choose very different query plans based on different estimates of index selectivity.

Whatever value of `innodb_stats_transient_sample_pages` works best for a system, set the option and leave it at that value. Choose a value that results in reasonably accurate estimates for all tables in your database without requiring excessive I/O. Because the statistics are automatically recalculated at various times other than on execution of `ANALYZE TABLE`, it does not make sense to increase the index sample size, run `ANALYZE TABLE`, then decrease sample size again.

Smaller tables generally require fewer index samples than larger tables. If your database has many large tables, consider using a higher value for `innodb_stats_transient_sample_pages` than if you have mostly smaller tables.

15.6.11.3 Estimating ANALYZE TABLE Complexity for InnoDB Tables

`ANALYZE TABLE` complexity for `InnoDB` tables is dependent on:

- The number of pages sampled, as defined by `innodb_stats_persistent_sample_pages`.
- The number of indexed columns in a table
- The number of partitions. If a table has no partitions, the number of partitions is considered to be 1.

Using these parameters, an approximate formula for estimating `ANALYZE TABLE` complexity would be:

The value of `innodb_stats_persistent_sample_pages` * number of indexed columns in a table * the number of partitions

Typically, the greater the resulting value, the greater the execution time for `ANALYZE TABLE`.

Note

`innodb_stats_persistent_sample_pages` defines the number of pages sampled at a global level. To set the number of pages sampled for an individual table, use the `STATS_SAMPLE_PAGES` option with `CREATE TABLE` or `ALTER TABLE`. For more information, see Section 15.6.11.1, "Configuring Persistent Optimizer Statistics Parameters".

If `innodb_stats_persistent=OFF`, the number of pages sampled is defined by `innodb_stats_transient_sample_pages`. See Section 15.6.11.2,

> "Configuring Non-Persistent Optimizer Statistics Parameters" for additional information.

For a more in-depth approach to estimating `ANALYZE TABLE` complexity, consider the following example.

In Big O notation, `ANALYZE TABLE` complexity is described as:

```
O(n_sample
  * (n_cols_in_uniq_i
     + n_cols_in_non_uniq_i
     + n_cols_in_pk * (1 + n_non_uniq_i))
  * n_part)
```

where:

- `n_sample` is the number of pages sampled (defined by `innodb_stats_persistent_sample_pages`)
- `n_cols_in_uniq_i` is total number of all columns in all unique indexes (not counting the primary key columns)
- `n_cols_in_non_uniq_i` is the total number of all columns in all non-unique indexes
- `n_cols_in_pk` is the number of columns in the primary key (if a primary key is not defined, `InnoDB` creates a single column primary key internally)
- `n_non_uniq_i` is the number of non-unique indexes in the table
- `n_part` is the number of partitions. If no partitions are defined, the table is considered to be a single partition.

Now, consider the following table (table `t`), which has a primary key (2 columns), a unique index (2 columns), and two non-unique indexes (two columns each):

```
CREATE TABLE t (
  a INT,
  b INT,
  c INT,
  d INT,
  e INT,
  f INT,
  g INT,
  h INT,
  PRIMARY KEY (a, b),
  UNIQUE KEY i1uniq (c, d),
  KEY i2nonuniq (e, f),
  KEY i3nonuniq (g, h)
);
```

For the column and index data required by the algorithm described above, query the `mysql.innodb_index_stats` persistent index statistics table for table `t`. The `n_diff_pfx%` statistics show the columns that are counted for each index. For example, columns `a` and `b` are counted for the primary key index. For the non-unique indexes, the primary key columns (a,b) are counted in addition to the user defined columns.

Note

For additional information about the `InnoDB` persistent statistics tables, see Section 15.6.11.1, "Configuring Persistent Optimizer Statistics Parameters"

```
  SELECT index_name, stat_name, stat_description
  FROM mysql.innodb_index_stats
  WHERE
  database_name='test' AND
  table_name='t' AND
```

```
stat_name like 'n_diff_pfx%';

+------------+--------------+------------------+
| index_name | stat_name    | stat_description |
+------------+--------------+------------------+
| PRIMARY    | n_diff_pfx01 | a                |
| PRIMARY    | n_diff_pfx02 | a,b              |
| i1uniq     | n_diff_pfx01 | c                |
| i1uniq     | n_diff_pfx02 | c,d              |
| i2nonuniq  | n_diff_pfx01 | e                |
| i2nonuniq  | n_diff_pfx02 | e,f              |
| i2nonuniq  | n_diff_pfx03 | e,f,a            |
| i2nonuniq  | n_diff_pfx04 | e,f,a,b          |
| i3nonuniq  | n_diff_pfx01 | g                |
| i3nonuniq  | n_diff_pfx02 | g,h              |
| i3nonuniq  | n_diff_pfx03 | g,h,a            |
| i3nonuniq  | n_diff_pfx04 | g,h,a,b          |
+------------+--------------+------------------+
```

Based on the index statistics data shown above and the table definition, the following values can be determined:

- `n_cols_in_uniq_i`, the total number of all columns in all unique indexes not counting the primary key columns, is 2 (`c` and `d`)
- `n_cols_in_non_uniq_i`, the total number of all columns in all non-unique indexes, is 4 (`e`, `f`, `g` and `h`)
- `n_cols_in_pk`, the number of columns in the primary key, is 2 (`a` and `b`)
- `n_non_uniq_i`, the number of non-unique indexes in the table, is 2 (`i2nonuniq` and `i3nonuniq`))
- `n_part`, the number of partitions, is 1.

You can now calculate `innodb_stats_persistent_sample_pages` * (2 + 4 + 2 * (1 + 2)) * 1 to determine the number of leaf pages that are scanned. With `innodb_stats_persistent_sample_pages` set to the default value of `20`, and with a default page size of 16 `KiB` (`innodb_page_size`=16384), you can then estimate that 20 * 12 * 16384 `bytes` are read for table `t`, or about 4 `MiB`.

Note

All 4 `MiB` may not be read from disk, as some leaf pages may already be cached in the buffer pool.

15.6.12 Configuring the Merge Threshold for Index Pages

You can configure the `MERGE_THRESHOLD` value for index pages. If the “page-full” percentage for an index page falls below the `MERGE_THRESHOLD` value when a row is deleted or when a row is shortened by an `UPDATE` operation, `InnoDB` attempts to merge the index page with a neighboring index page. The default `MERGE_THRESHOLD` value is 50, which is the previously hardcoded value. The minimum `MERGE_THRESHOLD` value is 1 and the maximum value is 50.

When the “page-full” percentage for an index page falls below 50%, which is the default `MERGE_THRESHOLD` setting, `InnoDB` attempts to merge the index page with a neighboring page. If both pages are close to 50% full, a page split can occur soon after the pages are merged. If this merge-split behavior occurs frequently, it can have an adverse affect on performance. To avoid frequent merge-splits, you can lower the `MERGE_THRESHOLD` value so that `InnoDB` attempts page merges at a lower “page-full” percentage. Merging pages at a lower page-full percentage leaves more room in index pages and helps reduce merge-split behavior.

The `MERGE_THRESHOLD` for index pages can be defined for a table or for individual indexes. A `MERGE_THRESHOLD` value defined for an individual index takes priority over a `MERGE_THRESHOLD` value defined for the table. If undefined, the `MERGE_THRESHOLD` value defaults to 50.

Setting MERGE_THRESHOLD for a Table

You can set the MERGE_THRESHOLD value for a table using the *table_option* COMMENT clause of the CREATE TABLE statement. For example:

```
CREATE TABLE t1 (
   id INT,
  KEY id_index (id)
) COMMENT='MERGE_THRESHOLD=45';
```

You can also set the MERGE_THRESHOLD value for an existing table using the *table_option* COMMENT clause with ALTER TABLE:

```
CREATE TABLE t1 (
   id INT,
  KEY id_index (id)
);

ALTER TABLE t1 COMMENT='MERGE_THRESHOLD=40';
```

Setting MERGE_THRESHOLD for Individual Indexes

To set the MERGE_THRESHOLD value for an individual index, you can use the *index_option* COMMENT clause with CREATE TABLE, ALTER TABLE, or CREATE INDEX, as shown in the following examples:

- Setting MERGE_THRESHOLD for an individual index using CREATE TABLE:

```
CREATE TABLE t1 (
   id INT,
  KEY id_index (id) COMMENT 'MERGE_THRESHOLD=40'
);
```

- Setting MERGE_THRESHOLD for an individual index using ALTER TABLE:

```
CREATE TABLE t1 (
   id INT,
  KEY id_index (id)
);

ALTER TABLE t1 DROP KEY id_index;
ALTER TABLE t1 ADD KEY id_index (id) COMMENT 'MERGE_THRESHOLD=40';
```

- Setting MERGE_THRESHOLD for an individual index using CREATE INDEX:

```
CREATE TABLE t1 (id INT);
CREATE INDEX id_index ON t1 (id) COMMENT 'MERGE_THRESHOLD=40';
```

Note

You cannot modify the MERGE_THRESHOLD value at the index level for GEN_CLUST_INDEX, which is the clustered index created by InnoDB when an InnoDB table is created without a primary key or unique key index. You can only modify the MERGE_THRESHOLD value for GEN_CLUST_INDEX by setting MERGE_THRESHOLD for the table.

Querying the MERGE_THRESHOLD Value for an Index

The current MERGE_THRESHOLD value for an index can be obtained by querying the INNODB_SYS_INDEXES table. For example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_INDEXES WHERE NAME='id_index' \G
*************************** 1. row ***************************
       INDEX_ID: 91
           NAME: id_index
       TABLE_ID: 68
           TYPE: 0
```

```
       N_FIELDS: 1
        PAGE_NO: 4
          SPACE: 57
MERGE_THRESHOLD: 40
```

You can use `SHOW CREATE TABLE` to view the `MERGE_THRESHOLD` value for a table, if explicitly defined using the *`table_option`* `COMMENT` clause:

```
mysql> SHOW CREATE TABLE t2 \G
*************************** 1. row ***************************
       Table: t2
Create Table: CREATE TABLE `t2` (
  `id` int(11) DEFAULT NULL,
  KEY `id_index` (`id`) COMMENT 'MERGE_THRESHOLD=40'
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4
```

Note

A `MERGE_THRESHOLD` value defined at the index level takes priority over a `MERGE_THRESHOLD` value defined for the table. If undefined, `MERGE_THRESHOLD` defaults to 50% (`MERGE_THRESHOLD=50`, which is the previously hardcoded value.

Likewise, you can use `SHOW INDEX` to view the `MERGE_THRESHOLD` value for an index, if explicitly defined using the *`index_option`* `COMMENT` clause:

```
mysql> SHOW INDEX FROM t2 \G
*************************** 1. row ***************************
        Table: t2
   Non_unique: 1
     Key_name: id_index
 Seq_in_index: 1
  Column_name: id
    Collation: A
  Cardinality: 0
     Sub_part: NULL
       Packed: NULL
         Null: YES
   Index_type: BTREE
      Comment:
Index_comment: MERGE_THRESHOLD=40
```

Measuring the Effect of MERGE_THRESHOLD Settings

The `INNODB_METRICS` table provides two counters that can be used to measure the effect of a `MERGE_THRESHOLD` setting on index page merges.

```
mysql> SELECT NAME, COMMENT FROM INFORMATION_SCHEMA.INNODB_METRICS
WHERE NAME like '%index_page_merge%';
+-----------------------------+---------------------------------------+
| NAME                        | COMMENT                               |
+-----------------------------+---------------------------------------+
| index_page_merge_attempts   | Number of index page merge attempts   |
| index_page_merge_successful | Number of successful index page merges |
+-----------------------------+---------------------------------------+
```

When lowering the `MERGE_THRESHOLD` value, the objectives are:

- A smaller number of page merge attempts and successful page merges
- A similar number of page merge attempts and successful page merges

A `MERGE_THRESHOLD` setting that is too small could result in large data files due to an excessive amount of empty page space.

For information about using `INNODB_METRICS` counters, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

15.7 InnoDB Tablespaces

This section covers topics related to InnoDB tablespaces.

15.7.1 Resizing the InnoDB System Tablespace

This section describes how to increase or decrease the size of the InnoDB system tablespace.

Increasing the Size of the InnoDB System Tablespace

The easiest way to increase the size of the InnoDB system tablespace is to configure it from the beginning to be auto-extending. Specify the autoextend attribute for the last data file in the tablespace definition. Then InnoDB increases the size of that file automatically in 64MB increments when it runs out of space. The increment size can be changed by setting the value of the innodb_autoextend_increment system variable, which is measured in megabytes.

You can expand the system tablespace by a defined amount by adding another data file:

1. Shut down the MySQL server.
2. If the previous last data file is defined with the keyword autoextend, change its definition to use a fixed size, based on how large it has actually grown. Check the size of the data file, round it down to the closest multiple of 1024 × 1024 bytes (= 1MB), and specify this rounded size explicitly in innodb_data_file_path.
3. Add a new data file to the end of innodb_data_file_path, optionally making that file auto-extending. Only the last data file in the innodb_data_file_path can be specified as auto-extending.
4. Start the MySQL server again.

For example, this tablespace has just one auto-extending data file ibdata1:

```
innodb_data_home_dir =
innodb_data_file_path = /ibdata/ibdata1:10M:autoextend
```

Suppose that this data file, over time, has grown to 988MB. Here is the configuration line after modifying the original data file to use a fixed size and adding a new auto-extending data file:

```
innodb_data_home_dir =
innodb_data_file_path = /ibdata/ibdata1:988M;/disk2/ibdata2:50M:autoextend
```

When you add a new data file to the system tablespace configuration, make sure that the filename does not refer to an existing file. InnoDB creates and initializes the file when you restart the server.

Decreasing the Size of the InnoDB System Tablespace

You cannot remove a data file from the system tablespace. To decrease the system tablespace size, use this procedure:

1. Use mysqldump to dump all your InnoDB tables, including InnoDB tables located in the MySQL database.

```
mysql> select table_name from information_schema.tables where table_schema='mysql' and engine='InnoDB';
+---------------------------+
| table_name                |
+---------------------------+
| columns_priv              |
| db                        |
```

```
| engine_cost               |
| gtid_executed             |
| help_category             |
| help_keyword              |
| help_relation             |
| help_topic                |
| innodb_index_stats        |
| innodb_table_stats        |
| plugin                    |
| procs_priv                |
| proxies_priv              |
| server_cost               |
| servers                   |
| slave_master_info         |
| slave_relay_log_info      |
| slave_worker_info         |
| tables_priv               |
| time_zone                 |
| time_zone_leap_second     |
| time_zone_name            |
| time_zone_transition      |
| time_zone_transition_type |
| user                      |
+---------------------------+
```

2. Stop the server.
3. Remove all the existing tablespace files (`*.ibd`), including the `ibdata` and `ib_log` files. Do not forget to remove `*.ibd` files for tables located in the MySQL database.
4. Configure a new tablespace.
5. Restart the server.
6. Import the dump files.

Note

If your databases only use the `InnoDB` engine, it may be simpler to dump **all** databases, stop the server, remove all databases and `InnoDB` log files, restart the server, and import the dump files.

15.7.2 Changing the Number or Size of InnoDB Redo Log Files

To change the number or the size of your `InnoDB` redo log files, perform the following steps:

1. Stop the MySQL server and make sure that it shuts down without errors.
2. Edit `my.cnf` to change the log file configuration. To change the log file size, configure `innodb_log_file_size`. To increase the number of log files, configure `innodb_log_files_in_group`.
3. Start the MySQL server again.

If `InnoDB` detects that the `innodb_log_file_size` differs from the redo log file size, it writes a log checkpoint, closes and removes the old log files, creates new log files at the requested size, and opens the new log files.

15.7.3 Using Raw Disk Partitions for the System Tablespace

You can use raw disk partitions as data files in the `InnoDB` system tablespace. This technique enables nonbuffered I/O on Windows and on some Linux and Unix systems without file system overhead. Perform tests with and without raw partitions to verify whether this change actually improves performance on your system.

When you use a raw disk partition, ensure that the user ID that runs the MySQL server has read and write privileges for that partition. For example, if you run the server as the `mysql` user, the partition must be readable and writeable by `mysql`. If you run the server with the `--memlock` option, the server must be run as `root`, so the partition must be readable and writeable by `root`.

The procedures described below involve option file modification. For additional information, see Section 4.2.6, “Using Option Files”.

Allocating a Raw Disk Partition on Linux and Unix Systems

1. When you create a new data file, specify the keyword `newraw` immediately after the data file size for the `innodb_data_file_path` option. The partition must be at least as large as the size that you specify. Note that 1MB in `InnoDB` is 1024 × 1024 bytes, whereas 1MB in disk specifications usually means 1,000,000 bytes.

   ```
   [mysqld]
   innodb_data_home_dir=
   innodb_data_file_path=/dev/hdd1:3Gnewraw;/dev/hdd2:2Gnewraw
   ```

2. Restart the server. `InnoDB` notices the `newraw` keyword and initializes the new partition. However, do not create or change any `InnoDB` tables yet. Otherwise, when you next restart the server, `InnoDB` reinitializes the partition and your changes are lost. (As a safety measure `InnoDB` prevents users from modifying data when any partition with `newraw` is specified.)

3. After `InnoDB` has initialized the new partition, stop the server, change `newraw` in the data file specification to `raw`:

   ```
   [mysqld]
   innodb_data_home_dir=
   innodb_data_file_path=/dev/hdd1:3Graw;/dev/hdd2:2Graw
   ```

4. Restart the server. `InnoDB` now permits changes to be made.

Allocating a Raw Disk Partition on Windows

On Windows systems, the same steps and accompanying guidelines described for Linux and Unix systems apply except that the `innodb_data_file_path` setting differs slightly on Windows.

1. When you create a new data file, specify the keyword `newraw` immediately after the data file size for the `innodb_data_file_path` option:

   ```
   [mysqld]
   innodb_data_home_dir=
   innodb_data_file_path=//./D::10Gnewraw
   ```

   The `//./` corresponds to the Windows syntax of `\\.\` for accessing physical drives. In the example above, `D:` is the drive letter of the partition.

2. Restart the server. `InnoDB` notices the `newraw` keyword and initializes the new partition.

3. After `InnoDB` has initialized the new partition, stop the server, change `newraw` in the data file specification to `raw`:

   ```
   [mysqld]
   innodb_data_home_dir=
   innodb_data_file_path=//./D::10Graw
   ```

4. Restart the server. `InnoDB` now permits changes to be made.

15.7.4 InnoDB File-Per-Table Tablespaces

Historically, all InnoDB tables and indexes were stored in the system tablespace. This monolithic approach was targeted at machines dedicated entirely to database processing, with carefully planned data growth, where any disk storage allocated to MySQL would never be needed for other purposes. InnoDB's file-per-table tablespace feature provides a more flexible alternative, where each InnoDB table and its indexes are stored in a separate .ibd data file. Each such .ibd data file represents an individual tablespace. This feature is controlled by the innodb_file_per_table configuration option, which is enabled by default.

Advantages of File-Per-Table Tablespaces

- You can reclaim disk space when truncating or dropping a table stored in a file-per-table tablepace. Truncating or dropping tables stored in the shared system tablespace creates free space internally in the system tablespace data files (ibdata files) which can only be used for new InnoDB data.

  Similarly, a table-copying ALTER TABLE operation on table that resides in a shared tablespace can increase the amount of space used by the tablespace. Such operations may require as much additional space as the data in the table plus indexes. The additional space required for the table-copying ALTER TABLE operation is not released back to the operating system as it is for file-per-table tablespaces.

- The TRUNCATE TABLE operation is faster when run on tables stored in file-per-table tablepaces.

- You can store specific tables on separate storage devices, for I/O optimization, space management, or backup purposes by specifying the location of each table using the syntax CREATE TABLE ... DATA DIRECTORY = *absolute_path_to_directory*, as explained in Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory”.

- You can run OPTIMIZE TABLE to compact or recreate a file-per-table tablespace. When you run an OPTIMIZE TABLE, InnoDB creates a new .ibd file with a temporary name, using only the space required to store actual data. When the optimization is complete, InnoDB removes the old .ibd file and replaces it with the new one. If the previous .ibd file grew significantly but the actual data only accounted for a portion of its size, running OPTIMIZE TABLE can reclaim the unused space.

- You can move individual InnoDB tables rather than entire databases.

- You can copy individual InnoDB tables from one MySQL instance to another (known as the transportable tablespace feature).

- Tables created in file-per-table tablespaces support features associated with compressed and dynamic row formats.

- You can enable more efficient storage for tables with large BLOB or TEXT columns using the dynamic row format.

- File-per-table tablespaces may improve chances for a successful recovery and save time when a corruption occurs, when a server cannot be restarted, or when backup and binary logs are unavailable.

- File-per-table tablespaces are convenient for per-table status reporting when copying or backing up tables.

- You can monitor table size at a file system level, without accessing MySQL.

- Common Linux file systems do not permit concurrent writes to a single file when innodb_flush_method is set to O_DIRECT. As a result, there are possible performance improvements when using file-per-table tablespaces in conjunction with innodb_flush_method.

- The system tablespace stores the data dictionary and undo logs, and is limited in size by InnoDB tablespace size limits. See Section 15.8.1.7, “Limits on InnoDB Tables”. With file-per-table tablespaces, each table has its own tablespace, which provides room for growth.

Potential Disadvantages of File-Per-Table Tablespaces

- With file-per-table tablespaces, each table may have unused space, which can only be utilized by rows of the same table. This could lead to wasted space if not properly managed.
- `fsync` operations must run on each open table rather than on a single file. Because there is a separate `fsync` operation for each file, write operations on multiple tables cannot be combined into a single I/O operation. This may require `InnoDB` to perform a higher total number of `fsync` operations.
- `mysqld` must keep one open file handle per table, which may impact performance if you have numerous tables in file-per-table tablespaces.
- More file descriptors are used.
- `innodb_file_per_table` is enabled by default. You may consider disabling it if backward compatibility with MySQL 5.5 or earlier is a concern. Disabling `innodb_file_per_table` prevents `ALTER TABLE` from moving an `InnoDB` table from the system tablespace to an individual `.ibd` file in cases where `ALTER TABLE` recreates the table (`ALGORITHM=COPY`).

  For example, when restructuring the clustered index for an `InnoDB` table, the table is re-created using the current setting for `innodb_file_per_table`. This behavior does not apply when adding or dropping `InnoDB` secondary indexes. When a secondary index is created without rebuilding the table, the index is stored in the same file as the table data, regardless of the current `innodb_file_per_table` setting. This behavior also does not apply to tables added to the system tablespace using `CREATE TABLE ... TABLESPACE` or `ALTER TABLE ... TABLESPACE` syntax. These tables are not affected by the `innodb_file_per_table` setting.
- If many tables are growing there is potential for more fragmentation which can impede `DROP TABLE` and table scan performance. However, when fragmentation is managed, having files in their own tablespace can improve performance.
- The buffer pool is scanned when dropping a file-per-table tablespace, which can take several seconds for buffer pools that are tens of gigabytes in size. The scan is performed with a broad internal lock, which may delay other operations. Tables in the system tablespace are not affected.
- The `innodb_autoextend_increment` variable, which defines increment size (in MB) for extending the size of an auto-extending shared tablespace file when it becomes full, does not apply to file-per-table tablespace files, which are auto-extending regardless of the `innodb_autoextend_increment` setting. The initial extensions are by small amounts, after which extensions occur in increments of 4MB.

15.7.4.1 Enabling and Disabling File-Per-Table Tablespaces

The `innodb_file_per_table` option is enabled by default.

To set the `innodb_file_per_table` option at startup, start the server with the `--innodb_file_per_table` command-line option, or add this line to the `[mysqld]` section of `my.cnf`:

```
[mysqld]
innodb_file_per_table=1
```

You can also set `innodb_file_per_table` dynamically, while the server is running:

```
SET GLOBAL innodb_file_per_table=1;
```

With `innodb_file_per_table` enabled, you can store `InnoDB` tables in a *`tbl_name`*`.ibd` file. Unlike the `MyISAM` storage engine, with its separate *`tbl_name`*`.MYD` and *`tbl_name`*`.MYI` files for indexes and data, `InnoDB` stores the data and the indexes together in a single `.ibd` file.

If you disable `innodb_file_per_table` in your startup options and restart the server, or disable it with the `SET GLOBAL` command, `InnoDB` creates new tables inside the system tablespace unless you have explicitly placed the table in file-per-table tablespace or general tablespace using the `CREATE TABLE ... TABLESPACE` option.

You can always read and write any `InnoDB` tables, regardless of the file-per-table setting.

To move a table from the system tablespace to its own tablespace, change the `innodb_file_per_table` setting and rebuild the table:

```
SET GLOBAL innodb_file_per_table=1;
ALTER TABLE table_name ENGINE=InnoDB;
```

Tables added to the system tablespace using `CREATE TABLE ... TABLESPACE` or `ALTER TABLE ... TABLESPACE` syntax are not affected by the `innodb_file_per_table` setting. To move these tables from the system tablespace to a file-per-table tablespace, they must be moved explicitly using `ALTER TABLE ... TABLESPACE` syntax.

Note

`InnoDB` always needs the system tablespace because it puts its internal data dictionary and undo logs there. The `.ibd` files are not sufficient for `InnoDB` to operate.

When a table is moved out of the system tablespace into its own `.ibd` file, the data files that make up the system tablespace remain the same size. The space formerly occupied by the table can be reused for new `InnoDB` data, but is not reclaimed for use by the operating system. When moving large `InnoDB` tables out of the system tablespace, where disk space is limited, you may prefer to enable `innodb_file_per_table` and recreate the entire instance using the `mysqldump` command. As mentioned above, tables added to the system tablespace using `CREATE TABLE ... TABLESPACE` or `ALTER TABLE ... TABLESPACE` syntax are not affected by the `innodb_file_per_table` setting. These tables must be moved individually.

15.7.5 Creating File-Per-Table Tablespaces Outside the Data Directory

To create a new `InnoDB` file-per-table tablespace in a specific location outside the MySQL data directory, use the `DATA DIRECTORY = absolute_path_to_directory` clause of the `CREATE TABLE` statement.

Plan the location in advance, because you cannot use the `DATA DIRECTORY` clause with the `ALTER TABLE` statement. The directory you specify could be on another storage device with particular performance or capacity characteristics, such as a fast SSD or a high-capacity HDD.

Within the destination directory, MySQL creates a subdirectory corresponding to the database name, and within that a .ibd file for the new table.

The following example demonstrates creating a file-per-table tablespace outside the MySQL data directory. It shows the `.ibd` created in the specified directory.

```
mysql> USE test;
Database changed

mysql> SHOW VARIABLES LIKE 'innodb_file_per_table';
+-----------------------+-------+
| Variable_name         | Value |
+-----------------------+-------+
| innodb_file_per_table | ON    |
+-----------------------+-------+
1 row in set (0.00 sec)

mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) DATA DIRECTORY = '/alternative/directory';
```

```
Query OK, 0 rows affected (0.03 sec)

# MySQL creates a .ibd file for the new table in a subdirectory that corresponding
# to the database name

db_user@ubuntu:~/alternative/directory/test$ ls
t1.ibd
```

You can also use `CREATE TABLE ... TABLESPACE` in combination with the `DATA DIRECTORY` clause to create a file-per-table tablespace outside the MySQL data directory. To do so, you must specify `innodb_file_per_table` as the tablespace name.

```
CREATE TABLE t2 (c1 INT PRIMARY KEY) TABLESPACE = innodb_file_per_table
  DATA DIRECTORY = '/alternative/directory';
```

You do not have to enable `innodb_file_per_table` when using this method.

Usage Notes:

- MySQL initially holds the `.ibd` file open, preventing you from dismounting the device, but might eventually close the table if the server is busy. Be careful not to accidentally dismount an external device while MySQL is running, or to start MySQL while the device is disconnected. Attempting to access a table when the associated `.ibd` file is missing causes a serious error that requires a server restart.

  A server restart issues errors and warnings if the `.ibd` file is not at the expected path. In this case, you can restore the tablespace `.ibd` file from a backup or drop the table to remove the information about it from the data dictionary.

- Before tables on an NFS-mounted volume, review potential issues outlined in Using NFS with MySQL.

- If you use an LVM snapshot, file copy, or other file-based mechanism to back up the `.ibd` file, always use the `FLUSH TABLES ... FOR EXPORT` statement first to make sure all changes that were buffered in memory are flushed to disk before the backup occurs.

- The `DATA DIRECTORY` clause is a supported alternative to using symbolic links, which has always been problematic and was never supported for individual `InnoDB` tables.

15.7.6 Copying File-Per-Table Tablespaces to Another Instance

This section describes how to copy a file-per-table tablespaces from one MySQL instance to another, otherwise known as the Transportable Tablespaces feature. This feature also supports partitioned `InnoDB` tables and individual `InnoDB` table partitions and subpartitions.

For information about other `InnoDB` table copying methods, see Section 15.8.1.3, “Moving or Copying InnoDB Tables”.

There are many reasons why you might copy an `InnoDB` file-per-table tablespace to a different instance:

- To run reports without putting extra load on a production server.

- To set up identical data for a table on a new slave server.

- To restore a backed-up version of a table or partition after a problem or mistake.

- As a faster way of moving data around than importing the results of a `mysqldump` command. The data is available immediately, rather than having to be re-inserted and the indexes rebuilt.

- To move a file-per-table tablespace to a server with storage medium that better suits system requirements. For example, you may want to have busy tables on an SSD device, or large tables on a high-capacity HDD device.

Limitations and Usage Notes

- The tablespace copy procedure is only possible when `innodb_file_per_table` is enabled, which is the default setting. Tables residing in the shared system tablespace cannot be quiesced.
- When a table is quiesced, only read-only transactions are allowed on the affected table.
- When importing a tablespace, the page size must match the page size of the importing instance.
- `ALTER TABLE ... DISCARD TABLESPACE` is supported for partitioned `InnoDB` tables, and `ALTER TABLE ... DISCARD PARTITION ... TABLESPACE` is supported for `InnoDB` table partitions.
- `DISCARD TABLESPACE` is not supported for tablespaces with a parent-child (primary key-foreign key) relationship when `foreign_key_checks` is set to `1`. Before discarding a tablespace for parent-child tables, set `foreign_key_checks=0`. Partitioned `InnoDB` tables do not support foreign keys.
- `ALTER TABLE ... IMPORT TABLESPACE` does not enforce foreign key constraints on imported data. If there are foreign key constraints between tables, all tables should be exported at the same (logical) point in time. Partitioned `InnoDB` tables do not support foreign keys.
- `ALTER TABLE ... IMPORT TABLESPACE` and `ALTER TABLE ... IMPORT PARTITION ... TABLESPACE` do not require a `.cfg` metadata file to import a tablespace. However, metadata checks are not performed when importing without a `.cfg` file, and a warning similar to the following is issued:

  ```
  Message: InnoDB: IO Read error: (2, No such file or directory) Error opening '.\
  test\t.cfg', will attempt to import without schema verification
  1 row in set (0.00 sec)
  ```

  The ability to import without a `.cfg` file may be more convenient when no schema mismatches are expected. Additionally, the ability to import without a `.cfg` file could be useful in crash recovery scenarios in which metadata cannot be collected from an `.ibd` file.

  If no `.cfg` file is used, `InnoDB` uses the equivalent of a `SELECT MAX(ai_col) FROM table_name FOR UPDATE` statement to initialize the in-memory auto-increment counter that is used in assigning values for to an `AUTO_INCREMENT` column. Otherwise, the current maximum auto-increment counter value is read from the `.cfg` metadata file. For related information, see InnoDB AUTO_INCREMENT Counter Initialization.
- Due to a `.cfg` metadata file limitation, schema mismatches are not reported for partition type or partition definition differences when importing tablespace files for partitioned tables. Column differences are reported.
- When running `ALTER TABLE ... DISCARD PARTITION ... TABLESPACE` and `ALTER TABLE ... IMPORT PARTITION ... TABLESPACE` on subpartitioned tables, both partition and subpartition table names are allowed. When a partition name is specified, subpartitions of that partition are included in the operation.
- Importing a tablespace file from another MySQL server instance works if both instances have GA (General Availability) status and their versions are within the same series. Otherwise, the file must have been created on the same server instance into which it is imported.
- In replication scenarios, `innodb_file_per_table` must be set to `ON` on both the master and slave.
- On Windows, `InnoDB` stores database, tablespace, and table names internally in lowercase. To avoid import problems on case-sensitive operating systems such as Linux and UNIX, create all databases, tablespaces, and tables using lowercase names. A convenient way to accomplish this is to add the following line to the `[mysqld]` section of your `my.cnf` or `my.ini` file before creating databases, tablespaces, or tables:

```
[mysqld]
lower_case_table_names=1
```

- `ALTER TABLE ... DISCARD TABLESPACE` and `ALTER TABLE ...IMPORT TABLESPACE` are not supported with tables that belong to an `InnoDB` general tablespace. For more information, see `CREATE TABLESPACE`.
- The default row format for `InnoDB` tables is configurable using the `innodb_default_row_format` configuration option. Attempting to import a table that does not explicitly define a row format (`ROW_FORMAT`), or that uses `ROW_FORMAT=DEFAULT`, could result in a schema mismatch error if the `innodb_default_row_format` setting on the source instance differs from the setting on the destination instance. For related information, see Section 15.10.2, “Specifying the Row Format for a Table”.
- When exporting a tablespace that is encrypted using the `InnoDB` tablespace encryption feature, `InnoDB` generates a `.cfp` file in addition to a `.cfg` metadata file. The `.cfp` file must be copied to the destination instance together with the `.cfg` file and tablespace file before performing the `ALTER TABLE ... IMPORT TABLESPACE` operation on the destination instance. The `.cfp` file contains a transfer key and an encrypted tablespace key. On import, `InnoDB` uses the transfer key to decrypt the tablespace key. For related information, see Section 15.7.10, “InnoDB Tablespace Encryption”.

15.7.6.1 Transportable Tablespace Examples

Note

If you are transporting tables that are encrypted using the `InnoDB` tablespace encryption, see Limitations and Usage Notes before you begin for additional procedural information.

Example 1: Copying an InnoDB Table to Another Instance

This procedure demonstrates how to copy a regular `InnoDB` table from a running MySQL server instance to another running instance. The same procedure with minor adjustments can be used to perform a full table restore on the same instance.

1. On the source instance, create a table if one does not exist:

```
mysql> use test;
mysql> CREATE TABLE t(c1 INT) engine=InnoDB;
```

2. On the destination instance, create a table if one does not exist:

```
mysql> use test;
mysql> CREATE TABLE t(c1 INT) engine=InnoDB;
```

3. On the destination instance, discard the existing tablespace. (Before a tablespace can be imported, `InnoDB` must discard the tablespace that is attached to the receiving table.)

```
mysql> ALTER TABLE t DISCARD TABLESPACE;
```

4. On the source instance, run `FLUSH TABLES ... FOR EXPORT` to quiesce the table and create the `.cfg` metadata file:

```
mysql> use test;
mysql> FLUSH TABLES t FOR EXPORT;
```

The metadata (`.cfg`) is created in the `InnoDB` data directory.

Note

The `FLUSH TABLES ... FOR EXPORT` statement ensures that changes to the named table have been flushed to disk so that a binary table copy can be made while the instance is running. When `FLUSH TABLES ... FOR EXPORT` is run, `InnoDB` produces a `.cfg` file in the same database directory as the table. The `.cfg` file contains metadata used for schema verification when importing the tablespace file.

5. Copy the `.ibd` file and `.cfg` metadata file from the source instance to the destination instance. For example:

```
shell> scp /path/to/datadir/test/t.{ibd,cfg} destination-server:/path/to/datadir/test
```

Note

The `.ibd` file and `.cfg` file must be copied before releasing the shared locks, as described in the next step.

6. On the source instance, use `UNLOCK TABLES` to release the locks acquired by `FLUSH TABLES ... FOR EXPORT`:

```
mysql> use test;
mysql> UNLOCK TABLES;
```

7. On the destination instance, import the tablespace:

```
mysql> use test;
mysql> ALTER TABLE t IMPORT TABLESPACE;
```

Note

The `ALTER TABLE ... IMPORT TABLESPACE` feature does not enforce foreign key constraints on imported data. If there are foreign key constraints between tables, all tables should be exported at the same (logical) point in time. In this case you would stop updating the tables, commit all transactions, acquire shared locks on the tables, and then perform the export operation.

Example 2: Copying an InnoDB Partitioned Table to Another Instance

This procedure demonstrates how to copy a partitioned `InnoDB` table from a running MySQL server instance to another running instance. The same procedure with minor adjustments can be used to perform a full restore of a partitioned `InnoDB` table on the same instance.

1. On the source instance, create a partitioned table if one does not exist. In the following example, a table with three partitions (p0, p1, p2) is created:

```
mysql> use test;
mysql> CREATE TABLE t1 (i int) ENGINE = InnoDB PARTITION BY KEY (i) PARTITIONS 3;
```

In the `/datadir/test` directory, there is a separate tablespace (`.ibd`) file for each of the three partitions.

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd
```

2. On the destination instance, create the same partitioned table:

```
mysql> use test;
mysql> CREATE TABLE t1 (i int) ENGINE = InnoDB PARTITION BY KEY (i) PARTITIONS 3;
```

In the */datadir/test* directory, there is a separate tablespace (.ibd) file for each of the three partitions.

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd
```

3. On the destination instance, discard the tablespace for the partitioned table. (Before the tablespace can be imported on the destination instance, the tablespace that is attached to the receiving table must be discarded.)

```
mysql> ALTER TABLE t1 DISCARD TABLESPACE;
```

The three .ibd files that make up the tablespace for the partitioned table are discarded from the */datadir/test* directory.

4. On the source instance, run FLUSH TABLES ... FOR EXPORT to quiesce the partitioned table and create the .cfg metadata files:

```
mysql> use test;
mysql> FLUSH TABLES t1 FOR EXPORT;
```

Metadata (.cfg) files, one for each tablespace (.ibd) file, are created in the */datadir/test* directory on the source instance:

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd
t1#P#p0.cfg  t1#P#p1.cfg  t1#P#p2.cfg
```

Note

FLUSH TABLES ... FOR EXPORT statement ensures that changes to the named table have been flushed to disk so that binary table copy can be made while the instance is running. When FLUSH TABLES ... FOR EXPORT is run, InnoDB produces a .cfg metadata file for the table's tablespace files in the same database directory as the table. The .cfg files contain metadata used for schema verification when importing tablespace files. FLUSH TABLES ... FOR EXPORT can only be run on the table, not on individual table partitions.

5. Copy the .ibd and .cfg files from the source instance database directory to the destination instance database directory. For example:

```
shell> scp /path/to/datadir/test/t1*.{ibd,cfg} destination-server:/path/to/datadir/test
```

Note

The .ibd and .cfg files must be copied before releasing the shared locks, as described in the next step.

6. On the source instance, use UNLOCK TABLES to release the locks acquired by FLUSH TABLES ... FOR EXPORT:

```
mysql> use test;
mysql> UNLOCK TABLES;
```

7. On the destination instance, import the tablespace for the partitioned table:

```
mysql> use test;
mysql> ALTER TABLE t1 IMPORT TABLESPACE;
```

Example 3: Copying InnoDB Table Partitions to Another Instance

This procedure demonstrates how to copy `InnoDB` table partitions from a running MySQL server instance to another running instance. The same procedure with minor adjustments can be used to perform a restore of `InnoDB` table partitions on the same instance. In the following example, a partitioned table with four partitions (p0, p1, p2, p3) is created on the source instance. Two of the partitions (p2 and p3) are copied to the destination instance.

1. On the source instance, create a partitioned table if one does not exist. In the following example, a table with four partitions (p0, p1, p2, p3) is created:

```
mysql> use test;
mysql> CREATE TABLE t1 (i int) ENGINE = InnoDB PARTITION BY KEY (i) PARTITIONS 4;
```

   In the *`/datadir`*`/test` directory, there is a separate tablespace (`.ibd`) file for each of the four partitions.

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd t1#P#p3.ibd
```

2. On the destination instance, create the same partitioned table:

```
mysql> use test;
mysql> CREATE TABLE t1 (i int) ENGINE = InnoDB PARTITION BY KEY (i) PARTITIONS 4;
```

   In the *`/datadir`*`/test` directory, there is a separate tablespace (`.ibd`) file for each of the four partitions.

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd t1#P#p3.ibd
```

3. On the destination instance, discard the tablespace partitions that you plan to import from the source instance. (Before tablespace partitions can be imported on the destination instance, the corresponding partitions that are attached to the receiving table must be discarded.)

```
mysql> ALTER TABLE t1 DISCARD PARTITION p2, p3 TABLESPACE;
```

   The `.ibd` files for the two discarded partitions are removed from the *`/datadir`*`/test` directory on the destination instance, leaving the following files:

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd
```

Note

When `ALTER TABLE ... DISCARD PARTITION ... TABLESPACE` is run on subpartitioned tables, both partition and subpartition table names are allowed. When a partition name is specified, subpartitions of that partition are included in the operation.

4. On the source instance, run `FLUSH TABLES ... FOR EXPORT` to quiesce the partitioned table and create the `.cfg` metadata files.

```
```

```
mysql> use test;
mysql> FLUSH TABLES t1 FOR EXPORT;
```

The metadata files (`.cfg` files) are created in the `/datadir/test` directory on the source instance. There is a `.cfg` file for each tablespace (`.ibd`) file.

```
mysql> \! ls /path/to/datadir/test/
t1#P#p0.ibd  t1#P#p1.ibd  t1#P#p2.ibd t1#P#p3.ibd
t1#P#p0.cfg  t1#P#p1.cfg  t1#P#p2.cfg t1#P#p3.cfg
```

Note

`FLUSH TABLES ... FOR EXPORT` statement ensures that changes to the named table have been flushed to disk so that binary table copy can be made while the instance is running. When `FLUSH TABLES ... FOR EXPORT` is run, `InnoDB` produces a `.cfg` metadata file for the table's tablespace files in the same database directory as the table. The `.cfg` files contain metadata used for schema verification when importing tablespace files. `FLUSH TABLES ... FOR EXPORT` can only be run on the table, not on individual table partitions.

5. Copy the `.ibd` and `.cfg` files from the source instance database directory to the destination instance database directory. In this example, only the `.ibd` and `.cfg` files for partition 2 (p2) and partition 3 (p3) are copied to the `data` directory on the destination instance. Partition 0 (p0) and partition 1 (p1) remain on the source instance.

```
shell> scp t1#P#p2.ibd  t1#P#p2.cfg t1#P#p3.ibd t1#P#p3.cfg destination-server:/path/to/datadir/test
```

Note

The `.ibd` files and `.cfg` files must be copied before releasing the shared locks, as described in the next step.

6. On the source instance, use `UNLOCK TABLES` to release the locks acquired by `FLUSH TABLES ... FOR EXPORT`:

```
mysql> use test;
mysql> UNLOCK TABLES;
```

7. On the destination instance, import the tablespace partitions (p2 and p3):

```
mysql> use test;
mysql> ALTER TABLE t1 IMPORT PARTITION p2, p3 TABLESPACE;
```

Note

When `ALTER TABLE ... IMPORT PARTITION ... TABLESPACE` is run on subpartitioned tables, both partition and subpartition table names are allowed. When a partition name is specified, subpartitions of that partition are included in the operation.

15.7.6.2 Transportable Tablespace Internals

The following information describes internals and error log messaging for the transportable tablespaces copy procedure for a regular `InnoDB` table.

When `ALTER TABLE ... DISCARD TABLESPACE` is run on the destination instance:

- The table is locked in X mode.

- The tablespace is detached from the table.

When `FLUSH TABLES ... FOR EXPORT` is run on the source instance:

- The table being flushed for export is locked in shared mode.
- The purge coordinator thread is stopped.
- Dirty pages are synchronized to disk.
- Table metadata is written to the binary `.cfg` file.

Expected error log messages for this operation:

```
2013-09-24T13:10:19.903526Z 2 [Note] InnoDB: Sync to disk of '"test"."t"' started.
2013-09-24T13:10:19.903586Z 2 [Note] InnoDB: Stopping purge
2013-09-24T13:10:19.903725Z 2 [Note] InnoDB: Writing table metadata to './test/t.cfg'
2013-09-24T13:10:19.904014Z 2 [Note] InnoDB: Table '"test"."t"' flushed to disk
```

When `UNLOCK TABLES` is run on the source instance:

- The binary .cfg file is deleted.
- The shared lock on the table or tables being imported is released and the purge coordinator thread is restarted.

Expected error log messages for this operation:

```
2013-09-24T13:10:21.181104Z 2 [Note] InnoDB: Deleting the meta-data file './test/t.cfg'
2013-09-24T13:10:21.181180Z 2 [Note] InnoDB: Resuming purge
```

When `ALTER TABLE ... IMPORT TABLESPACE` is run on the destination instance, the import algorithm performs the following operations for each tablespace being imported:

- Each tablespace page is checked for corruption.
- The space ID and log sequence numbers (LSNs) on each page are updated
- Flags are validated and LSN updated for the header page.
- Btree pages are updated.
- The page state is set to dirty so that it is written to disk.

Expected error log messages for this operation:

```
2013-07-18 15:15:01 34960 [Note] InnoDB: Importing tablespace for table 'test/t' that was exported from
2013-07-18 15:15:01 34960 [Note] InnoDB: Phase I - Update all pages
2013-07-18 15:15:01 34960 [Note] InnoDB: Sync to disk
2013-07-18 15:15:01 34960 [Note] InnoDB: Sync to disk - done!
2013-07-18 15:15:01 34960 [Note] InnoDB: Phase III - Flush changes to disk
2013-07-18 15:15:01 34960 [Note] InnoDB: Phase IV - Flush complete
```

Note

You may also receive a warning that a tablespace is discarded (if you discarded the tablespace for the destination table) and a message stating that statistics could not be calculated due to a missing `.ibd` file:

```
2013-07-18 15:14:38 34960 [Warning] InnoDB: Table "test"."t" tablespace is set as dis
2013-07-18 15:14:38 7f34d9a37700 InnoDB: cannot calculate statistics for table "test"
http://dev.mysql.com/doc/refman/5.8/en/innodb-troubleshooting.html
```

15.7.7 Storing InnoDB Undo Logs in Separate Tablespaces

Undo logs may be stored in one or more undo tablespaces instead of the system tablespace. The I/O patterns for undo logs make undo tablespaces good candidates for SSD storage, while keeping the system tablespace on hard disk storage. Undo tablespaces or individual segments inside those tablespaces cannot be dropped. However, undo logs stored in undo tablespaces can be truncated. For more information, see Section 15.7.8, “Truncating Undo Logs That Reside in Undo Tablespaces”.

Important

Prior to MySQL 8.0.2, the number of undo tablespaces may only be configured when initializing a new MySQL instance, as the `innodb_undo_tablespaces` option can only be set at initialization time. As of 8.0.2, `innodb_undo_tablespaces` may be set at startup or while the server is running.

To configure separate undo tabespaces for a MySQL instance, perform the following steps. It is assumed that you are performing the procedure on a test instance prior to deploying the configuration to a production system.

1. Specify a directory location for undo tablespaces using the `innodb_undo_directory` configuration option. If a directory location is not specified, undo tablespaces are created in the MySQL data directory.

2. Define the number of rollback segments using the `innodb_rollback_segments` configuration option. Start with a relatively low value and increase it incrementally over time to examine the effect on performance.

   - Prior to MySQL 8.0.2, one rollback segment is always assigned to the system tablespace. Therefore, to allocate rollback segments to undo tablespaces, set `innodb_rollback_segments` to a value greater than 1. For example, if you have two undo tablespaces, set `innodb_rollback_segments` to 3 to assign one rollback segment to each of the two undo tablespaces. Rollback segments are distributed among undo tablespaces in a circular fashion.

   - As of MySQL 8.0.2, `innodb_rollback_segments` defines the number of rollback segments allocated to each undo tablespace.

   When you configure separate undo tablespaces, rollback segments in the system tablespace are rendered inactive.

3. Define the number of undo tablespaces using the `innodb_undo_tablespaces` option.

   Prior to MySQL 8.0.2, the `innodb_undo_tablespaces` value can only be set when initializing the MySQL instance and cannot be changed afterward, so if you are uncertain about the optimal value, estimate on the high side. As of MySQL 8.0.2, `innodb_undo_tablespaces` can be configured at startup or while the server is running.

4. If you are using MySQL 8.0.1 or earlier, create a new MySQL test instance using the option values you have chosen.

5. Use a realistic workload on your test instance with data volume similar to your production servers to test the configuration.

6. Benchmark the performance of I/O intensive workloads.

7. Periodically increase the value of `innodb_rollback_segments` and rerun performance tests until there are no further improvements in I/O performance.

15.7.8 Truncating Undo Logs That Reside in Undo Tablespaces

To truncate undo logs that reside in undo tablespaces, the MySQL instance must be configured with a minimum of two undo tablespaces. A minimum of two undo tablespaces ensures that one undo tablespace remains active while the other is taken offline to be truncated. The number of undo tablespaces is defined by the `innodb_undo_tablespaces` option. The default value is 2 as of MySQL 8.0.2, and 0 before that. Use this statement to check the value of `innodb_undo_tablespaces`:

```
mysql> SELECT @@innodb_undo_tablespaces;
+--------------------------+
| @@innodb_undo_tablespaces |
+--------------------------+
|                        2 |
+--------------------------+
```

For information about configuring undo tablespaces, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

Enabling Truncation of Undo Tablespaces

To truncate undo logs that reside in undo tablespaces, enable `innodb_undo_log_truncate`.

```
mysql> SET GLOBAL innodb_undo_log_truncate=ON;
```

When `innodb_undo_log_truncate` is enabled, undo tablespace files that exceed the size limit defined by `innodb_max_undo_log_size` are marked for truncation. `innodb_max_undo_log_size` is a dynamic global variable with a default value of 1024 MiB (1073741824 bytes).

```
mysql> SELECT @@innodb_max_undo_log_size;
+----------------------------+
| @@innodb_max_undo_log_size |
+----------------------------+
|                 1073741824 |
+----------------------------+
```

You can configure `innodb_max_undo_log_size` using a `SET GLOBAL` statement:

```
mysql> SET GLOBAL innodb_max_undo_log_size=2147483648;
```

When `innodb_undo_log_truncate` is enabled:

1. Undo tablespaces that exceed the `innodb_max_undo_log_size` setting are marked for truncation. Selection of an undo tablespace for truncation is performed in a circular fashion to avoid truncating the same undo tablespace each time.
2. Rollback segments residing in the selected undo tablespace are made inactive so that they are not assigned to new transactions. Existing transactions that are currently using rollback segments are allowed to complete.
3. The purge system frees rollback segments that are no longer needed.
4. After all rollback segments in the undo tablespace are freed, the truncate operation runs and the undo tablespace is truncated to its initial size. The initial size of an undo tablespace file is 10MiB.

Note

The size of an undo tablespace after a truncate operation may be larger than 10MiB due to immediate use following the completion of the operation. The `innodb_undo_directory` option defines the location of undo tablespace files. The default value of “.” represents the directory where `InnoDB` creates other log files by default.

```
mysql> SELECT @@innodb_undo_directory;
+--------------------------+
| @@innodb_undo_directory |
+--------------------------+
| .                        |
+--------------------------+
```

5. The rollback segments are reactivated so that they can be assigned to new transactions.

Expediting Truncation of Undo Tablespace Files

An undo tablespace cannot be truncated until its rollback segments are freed. Normally, the purge system frees rollback segments once every 128 times that purge is invoked. To expedite the truncation of undo tablespaces, use the `innodb_purge_rseg_truncate_frequency` option to temporarily increase the frequency with which the purge system frees rollback segments. The default `innodb_purge_rseg_truncate_frequency` setting is 128, which is also the maximum value.

```
mysql> SELECT @@innodb_purge_rseg_truncate_frequency;
+-----------------------------------------+
| @@innodb_purge_rseg_truncate_frequency |
+-----------------------------------------+
|                                     128 |
+-----------------------------------------+
```

To increase the frequency with which the purge thread frees rollback segments, decrease the value of `innodb_purge_rseg_truncate_frequency`. For example:

```
mysql> SET GLOBAL innodb_purge_rseg_truncate_frequency=32;
```

Performance Impact of Truncating Undo Tablespace Files Online

While an undo tablespace is truncated, rollback segments in that tablespace are temporarily deactivated. The remaining active rollback segments in the other undo tablespaces assume responsibility for the entire system load, which may result in a slight performance degradation. The degree of performance degradation depends on a number of factors including:

- Number of undo tablespaces
- Number of undo logs
- Undo tablespace size
- Speed of the I/O susbsystem
- Existing long running transactions
- System load

15.7.9 InnoDB General Tablespaces

A general tablespace is a shared `InnoDB` tablespace that is created using `CREATE TABLESPACE` syntax. General tablespace capabilities and features are described under the following topics in this section:

- General Tablespace Capabilities
- Creating a General Tablespace
- Adding Tables to a General Tablespace
- General Tablespace Row Format Support
- Moving Non-Partitioned Tables Between Tablespaces Using ALTER TABLE

- General Tablespace Table Partition Support
- Moving Table Partitions Between Tablespaces Using ALTER TABLE
- Dropping a General Tablespace
- General Tablespace Limitations

General Tablespace Capabilities

The general tablespace feature provides the following capabilities:

- Similar to the system tablespace, general tablespaces are shared tablespaces that can store data for multiple tables.
- General tablespaces have a potential memory advantage over file-per-table tablespaces. The server keeps tablespace metadata in memory for the lifetime of a tablespace. Multiple tables in fewer general tablespaces consume less memory for tablespace metadata than the same number of tables in separate file-per-table tablespaces.
- General tablespace data files may be placed in a directory relative to or independent of the MySQL data directory, which provides you with many of the data file and storage management capabilities of file-per-table tablespaces. As with file-per-table tablespaces, the ability to place data files outside of the MySQL data directory allows you to manage performance of critical tables separately, setup RAID or DRBD for specific tables, or bind tables to particular disks, for example.
- General tablespaces support both Antelope and Barracuda file formats, and therefore support all table row formats and associated features. With support for both file formats, general tablespaces have no dependence on `innodb_file_format` or `innodb_file_per_table` settings, nor do these variables have any effect on general tablespaces.
- The `TABLESPACE` option can be used with `CREATE TABLE` to create tables in a general tablespaces, file-per-table tablespace, or in the system tablespace.
- The `TABLESPACE` option can be used with `ALTER TABLE` to move tables between general tablespaces, file-per-table tablespaces, and the system tablespace. Previously, it was not possible to move a table from a file-per-table tablespace to the system tablespace. With the general tablespace feature, you can now do so.

Creating a General Tablespace

General tablespaces are created using `CREATE TABLESPACE` syntax.

```
CREATE TABLESPACE tablespace_name
    ADD DATAFILE 'file_name'
    [FILE_BLOCK_SIZE = value]
        [ENGINE [=] engine_name]
```

A general tablespace may be created in the MySQL data directory or in a directory outside of the MySQL data directory. To avoid conflicts with implicitly created file-per-table tablespaces, creating a general tablespace in a subdirectory under the MySQL data directory is not supported. Also, when creating a general tablespace outside of the MySQL data directory, the directory must exist prior to creating the tablespace.

Examples:

Creating a general tablespace in the MySQL data directory:

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
```

Creating a general tablespace in a directory outside of the MySQL data directory:

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE '/my/tablespace/directory/ts1.ibd' Engine=InnoDB;
```

You can specify a path that is relative to the MySQL data directory as long as the tablespace directory is not under the MySQL data directory. In this example, the `my_tablespace` directory is at the same level as the MySQL data directory:

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE '../my_tablespace/ts1.ibd' Engine=InnoDB;
```

Note

The `ENGINE = InnoDB` clause must be defined as part of the `CREATE TABLESPACE` statement or `InnoDB` must be defined as the default storage engine (`default_storage_engine=InnoDB`).

Adding Tables to a General Tablespace

After creating an `InnoDB` general tablespace, you can use `CREATE TABLE *tbl_name* ... TABLESPACE [=] *tablespace_name*` or `ALTER TABLE *tbl_name* TABLESPACE [=] *tablespace_name*` to add tables to the tablespace, as shown in the following examples:

`CREATE TABLE`:

```
mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=COMPACT;
```

`ALTER TABLE`:

```
mysql> ALTER TABLE t2 TABLESPACE ts1;
```

For detailed syntax information, see `CREATE TABLE` and `ALTER TABLE`.

General Tablespace Row Format Support

General tablespaces support all table row formats (`REDUNDANT`, `COMPACT`, `DYNAMIC`, `COMPRESSED`) with the caveat that compressed and uncompressed tables cannot coexist in the same general tablespace due to different physical page sizes.

For a general tablespace to contain compressed tables (`ROW_FORMAT=COMPRESSED`), `FILE_BLOCK_SIZE` must be specified, and the `FILE_BLOCK_SIZE` value must be a valid compressed page size in relation to the `innodb_page_size` value. Also, the physical page size of the compressed table (`KEY_BLOCK_SIZE`) must be equal to `FILE_BLOCK_SIZE/1024`. For example, if `innodb_page_size=16K` and `FILE_BLOCK_SIZE=8K`, the `KEY_BLOCK_SIZE` of the table must be 8.

The following table shows permitted `FILE_BLOCK_SIZE` and `KEY_BLOCK_SIZE` values for each `innodb_page_size` value. `FILE_BLOCK_SIZE` values may also be specified in bytes. To determine a valid `KEY_BLOCK_SIZE` value for a given `FILE_BLOCK_SIZE`, divide the `FILE_BLOCK_SIZE` value by 1024. Table compression is not support for 32K and 64K `InnoDB` page sizes. For more information about `KEY_BLOCK_SIZE`, see `CREATE TABLE`, and Section 15.9.1.2, “Creating Compressed Tables”.

Table 15.5 FILE_BLOCK_SIZE and KEY_BLOCK_SIZE Values for CREATE TABLESPACE

InnoDB Page Size (innodb_page_size)	Permitted FILE_BLOCK_SIZE Values	Permitted KEY_BLOCK_SIZE Values
64K	64K (65536)	Compression is not supported
32K	32K (32768)	Compression is not supported
16K	16K (16384)	N/A: If `innodb_page_size` is equal to `FILE_BLOCK_SIZE`, the tablespace cannot contain a compressed table.
	8K (8192)	8
	4K (4096)	4
	2K (2048)	2

InnoDB Page Size (innodb_page_size)	Permitted FILE_BLOCK_SIZE Values	Permitted KEY_BLOCK_SIZE Values
	1K (1024)	1
8K	8K (8192)	N/A: If innodb_page_size is equal to FILE_BLOCK_SIZE, the tablespace cannot contain a compressed table.
	4K (4096)	4
	2K (2048)	2
	1K (1024)	1
4K	4K (4096)	N/A: If innodb_page_size is equal to FILE_BLOCK_SIZE, the tablespace cannot contain a compressed table.
	2K (2048)	2
	1K (1024)	1

This example demonstrates creating a general tablespace and adding a compressed table. The example assumes a default innodb_page_size of 16K. The FILE_BLOCK_SIZE of 8192 requires that the compressed table have a KEY_BLOCK_SIZE of 8.

```
mysql> CREATE TABLESPACE `ts2` ADD DATAFILE 'ts2.ibd' FILE_BLOCK_SIZE = 8192 Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t4 (c1 INT PRIMARY KEY) TABLESPACE ts2 ROW_FORMAT=COMPRESSED
KEY_BLOCK_SIZE=8;
Query OK, 0 rows affected (0.00 sec)
```

If you do not specify FILE_BLOCK_SIZE when creating a general tablespace, FILE_BLOCK_SIZE defaults to innodb_page_size. When FILE_BLOCK_SIZE is equal to innodb_page_size, the tablespace may only contain tables with an uncompressed row format (COMPACT, REDUNDANT, and DYNAMIC row formats).

Moving Non-Partitioned Tables Between Tablespaces Using ALTER TABLE

You can use ALTER TABLE with the TABLESPACE option to move a non-partitioned InnoDB table to an existing general tablespace, to a new file-per-table tablespace, or to the system tablespace.

To move a non-partitioned table from a file-per-table tablespace or from the system tablespace to a general tablespace, specify the name of the general tablespace. The general tablespace must exist. See CREATE TABLESPACE for more information.

```
ALTER TABLE tbl_name TABLESPACE [=] tablespace_name
```

To move a non-partitioned table from a general tablespace or file-per-table tablespace to the system tablespace, specify innodb_system as the tablespace name.

```
ALTER TABLE tbl_name ... TABLESPACE [=] innodb_system
```

To move a non-partitioned table from the system tablespace or a general tablespace to a file-per-table tablespace, specify innodb_file_per_table as the tablespace name.

```
ALTER TABLE tbl_name ... TABLESPACE [=] innodb_file_per_table
```

ALTER TABLE ... TABLESPACE operations always cause a full table rebuild, even if the TABLESPACE attribute has not changed from its previous value.

ALTER TABLE ... TABLESPACE syntax does not support moving a table from a temporary tablespace to a persistent tablespace.

The DATA DIRECTORY clause is permitted with CREATE TABLE ... TABLESPACE=innodb_file_per_table but is otherwise not supported for use in combination with the TABLESPACE option.

General Tablespace Table Partition Support

The TABLESPACE option may be used to assign individual table partitions or subpartitions to a general tablespace, a separate file-per-table tablespace, or the system tablespace. All partitions must belong to the same storage engine. Usage is demonstrated in the following examples.

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
mysql> CREATE TABLESPACE `ts2` ADD DATAFILE 'ts2.ibd' Engine=InnoDB;

mysql> CREATE TABLE t1 (a INT, b INT) ENGINE = InnoDB
    ->  PARTITION BY RANGE(a) SUBPARTITION BY KEY(b) (
    ->    PARTITION p1 VALUES LESS THAN (100) TABLESPACE=`ts1`,
    ->    PARTITION p2 VALUES LESS THAN (1000) TABLESPACE=`ts2`,
    ->    PARTITION p3 VALUES LESS THAN (10000) TABLESPACE `innodb_file_per_table`,
    ->    PARTITION p4 VALUES LESS THAN (100000) TABLESPACE `innodb_system`);

mysql> CREATE TABLE t2 (a INT, b INT) ENGINE = InnoDB
    ->  PARTITION BY RANGE(a) SUBPARTITION BY KEY(b) (
    ->    PARTITION p1 VALUES LESS THAN (100) TABLESPACE=`ts1`
    ->      (SUBPARTITION sp1,
    ->       SUBPARTITION sp2),
    ->    PARTITION p2 VALUES LESS THAN (1000)
    ->      (SUBPARTITION sp3,
    ->       SUBPARTITION sp4 TABLESPACE=`ts2`),
    ->    PARTITION p3 VALUES LESS THAN (10000)
    ->      (SUBPARTITION sp5 TABLESPACE `innodb_system`,
    ->       SUBPARTITION sp6 TABLESPACE `innodb_file_per_table`));
```

The TABLESPACE option is also supported with ALTER TABLE.

```
mysql> ALTER TABLE t1 ADD PARTITION (PARTITION p5 VALUES LESS THAN (1000000) TABLESPACE = `ts1`);
```

Note

If the TABLESPACE = *tablespace_name* option is not defined, the ALTER TABLE ... ADD PARTITION operation adds the partition to the table's default tablespace, which can be specified at the table level during CREATE TABLE or ALTER TABLE.

An ALTER TABLE *tbl_name* TABLESPACE [=] *tablespace_name* operation on a *partitioned* table only modifies the table's default tablespace. It does not move the table partitions. However, after changing the default tablespace, an operation that rebuilds the table, such as an ALTER TABLE operation that uses ALGORITHM=COPY, moves the partitions to the default tablespace if another tablespace is not defined explicitly using the TABLESPACE [=] *tablespace_name* clause.

To verify that partitions were placed in the specified tablespaces, you can query INFORMATION_SCHEMA.INNODB_SYS_TABLES:

```
mysql> SELECT NAME, SPACE, SPACE_TYPE FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES
    -> WHERE NAME LIKE '%t1%';
+----------------------+-------+------------+
| NAME                 | SPACE | SPACE_TYPE |
+----------------------+-------+------------+
| test/t1#P#p1#SP#p1sp0 |    57 | General    |
| test/t1#P#p2#SP#p2sp0 |    58 | General    |
| test/t1#P#p3#SP#p3sp0 |    59 | Single     |
| test/t1#P#p4#SP#p4sp0 |     0 | System     |
| test/t1#P#p5#SP#p5sp0 |    57 | General    |
+----------------------+-------+------------+

mysql> SELECT NAME, SPACE, SPACE_TYPE FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES
```

```
    -> WHERE NAME LIKE '%t2%';
+--------------------+-------+------------+
| NAME               | SPACE | SPACE_TYPE |
+--------------------+-------+------------+
| test/t2#P#p1#SP#sp1 |    57 | General    |
| test/t2#P#p1#SP#sp2 |    57 | General    |
| test/t2#P#p2#SP#sp3 |    60 | Single     |
| test/t2#P#p2#SP#sp4 |    58 | General    |
| test/t2#P#p3#SP#sp5 |     0 | System     |
| test/t2#P#p3#SP#sp6 |    61 | Single     |
+--------------------+-------+------------+
```

Moving Table Partitions Between Tablespaces Using ALTER TABLE

To move table partitions to a different tablespace, you must move each partition using an `ALTER TABLE tbl_name REORGANIZE PARTITION` statement.

The following example demonstrates how to move table partitions to a different tablespace. `INFORMATION_SCHEMA.INNODB_SYS_TABLES` and `INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES` are queried to verify that partitions are placed in the expected tablespace.

Note

If the `TABLESPACE = tablespace_name` option is not defined in the `REORGANIZE PARTITION` statement, `InnoDB` moves the partition to the table's default tablespace. In the example that follows, tablespace `ts1`, which is defined at the table level, is the default tablespace for table `t1`. Partition `P3` is moved from the system tablespace to tablespace `ts1` since no `TABLESPACE` option is specified in the `ALTER TABLE t1 REORGANIZE PARTITION` statement for partition `P3`.

An operation that rebuilds the table, such as an `ALTER TABLE` operation that uses `ALGORITHM=COPY`, moves partitions to the default tablespace if partitions reside in a different tablespace that is not defined explicitly using the `TABLESPACE [=] tablespace_name` clause.

```
mysql> CREATE TABLESPACE ts1 ADD DATAFILE 'ts1.ibd';
mysql> CREATE TABLESPACE ts2 ADD DATAFILE 'ts2.ibd';

mysql> CREATE TABLE t1 ( a INT NOT NULL, PRIMARY KEY (a))
    ->  ENGINE=InnoDB TABLESPACE ts1
    ->  PARTITION BY RANGE (a) PARTITIONS 3 (
    ->    PARTITION P1 VALUES LESS THAN (2),
    ->    PARTITION P2 VALUES LESS THAN (4) TABLESPACE `innodb_file_per_table`,
    ->    PARTITION P3 VALUES LESS THAN (6) TABLESPACE `innodb_system`);


mysql> SELECT A.NAME as partition_name, A.SPACE_TYPE as space_type, B.NAME as space_name
    -> FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES A
    -> LEFT JOIN INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES B
    -> ON A.SPACE = B.SPACE WHERE A.NAME LIKE '%t1%' ORDER BY A.NAME;
+----------------+------------+--------------+
| partition_name | space_type | space_name   |
+----------------+------------+--------------+
| test/t1#P#P1   | General    | ts1          |
| test/t1#P#P2   | Single     | test/t1#P#P2 |
| test/t1#P#P3   | System     | NULL         |
+----------------+------------+--------------+

mysql> ALTER TABLE t1 REORGANIZE PARTITION P1
    -> INTO (PARTITION P1 VALUES LESS THAN (2) TABLESPACE = `ts2`);

mysql> ALTER TABLE t1 REORGANIZE PARTITION P2
    -> INTO (PARTITION P2 VALUES LESS THAN (4) TABLESPACE = `ts2`);

mysql> ALTER TABLE t1 REORGANIZE PARTITION P3
```

```
    -> INTO (PARTITION P3 VALUES LESS THAN (6));

mysql> SELECT A.NAME AS partition_name, A.SPACE_TYPE AS space_type, B.NAME AS space_name
    -> FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES A
    -> LEFT JOIN INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES B
    -> ON A.SPACE = B.SPACE WHERE A.NAME LIKE '%t1%' ORDER BY A.NAME;
+----------------+------------+------------+
| partition_name | space_type | space_name |
+----------------+------------+------------+
| test/t1#P#P1   | General    | ts2        |
| test/t1#P#P2   | General    | ts2        |
| test/t1#P#P3   | General    | ts1        |
+----------------+------------+------------+
```

Dropping a General Tablespace

The `DROP TABLESPACE` statement is used to drop an `InnoDB` general tablespace.

All tables must be dropped from the tablespace prior to a `DROP TABLESPACE` operation. If the tablespace is not empty, `DROP TABLESPACE` returns an error.

A general `InnoDB` tablespace is not deleted automatically when the last table in the tablespace is dropped. The tablespace must be dropped explicitly using `DROP TABLESPACE` *`tablespace_name`*.

A general tablespace does not belong to any particular database. A `DROP DATABASE` operation can drop tables that belong to a general tablespace but it cannot drop the tablespace, even if the `DROP DATABASE` operation drops all tables that belong to the tablespace. A general tablespace must be dropped explicitly using `DROP TABLESPACE` *`tablespace_name`*.

Similar to the system tablespace, truncating or dropping tables stored in a general tablespace creates free space internally in the general tablespace .ibd data file which can only be used for new `InnoDB` data. Space is not released back to the operating system as it is when a file-per-table tablespace is deleted during a `DROP TABLE` operation.

This example demonstrates how to drop an `InnoDB` general tablespace. The general tablespace `ts1` is created with a single table. The table must be dropped before dropping the tablespace.

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts10 Engine=InnoDB;
Query OK, 0 rows affected (0.02 sec)

mysql> DROP TABLE t1;
Query OK, 0 rows affected (0.01 sec)

mysql> DROP TABLESPACE ts1;
Query OK, 0 rows affected (0.01 sec)
```

Note

`tablespace_name` is a case-sensitive identifier in MySQL.

General Tablespace Limitations

- A generated or existing tablespace cannot be changed to a general tablespace.
- Creation of temporary general tablespaces is not supported.
- General tablespaces do not support temporary tables.
- Similar to the system tablespace, truncating or dropping tables stored in a general tablespace creates free space internally in the general tablespace .ibd data file which can only be used for new `InnoDB` data. Space is not released back to the operating system as it is for file-per-table tablespaces.

Additionally, a table-copying `ALTER TABLE` operation on table that resides in a shared tablespace (a general tablespace or the system tablespace) can increase the amount of space used by the tablespace. Such operations require as much additional space as the data in the table plus indexes. The additional space required for the table-copying `ALTER TABLE` operation is not released back to the operating system as it is for file-per-table tablespaces.

- `ALTER TABLE ... DISCARD TABLESPACE` and `ALTER TABLE ...IMPORT TABLESPACE` are not supported for tables that belong to a general tablespace.

For more information see Section 13.1.16, “CREATE TABLESPACE Syntax”.

15.7.10 InnoDB Tablespace Encryption

`InnoDB` supports data encryption for `InnoDB` tables stored in file-per-table tablespaces. This feature provides at-rest encryption for physical tablespace data files.

`InnoDB` tablespace encryption uses a two tier encryption key architecture, consisting of a master encryption key and tablespace keys. When an `InnoDB` table is encrypted, a tablespace key is encrypted and stored in the tablespace header. When an application or authenticated user wants to access encrypted tablespace data, `InnoDB` uses the master encryption key to decrypt the tablespace key. The master encryption key is stored in a `keyring` file in the location specified by the `keyring_file_data` configuration option. The decrypted version of a tablespace key never changes, but the master encryption key may be changed as required. This action is referred to as *master key rotation*.

The `InnoDB` tablespace encryption feature relies on a keyring plugin for master encryption key management.

`InnoDB` tablespace encryption supports the Advanced Encryption Standard (AES) block-based encryption algorithm. It uses Electronic Codebook (ECB) block encryption mode for tablespace key encryption and Cipher Block Chaining (CBC) block encryption mode for data encryption.

Note

The `InnoDB` tablespace encryption feature provided with MySQL Community Edition is not intended as a regulatory compliance solution. Security standards such as PCI, FIPS, and others require use of key management systems to secure, manage, and protect keys in key vaults or hardware security modules (HSMs).

For frequently asked questions about the `InnoDB` tablespace encryption feature, see Section A.15, “MySQL 8.0 FAQ: InnoDB Tablespace Encryption”.

InnoDB Tablespace Encryption Prerequisites

- The `keyring_file` plugin must be installed. Keyring plugin installation is performed at startup using the `--early-plugin-load` option. Early loading ensures that the plugin is available prior to initialization of the `InnoDB` storage engine. For keyring plugin installation and configuration instructions, see Section 6.5.4, “The MySQL Keyring”.

Important

Once encrypted tables are created in a MySQL instance, the keyring plugin must continue to be loaded using the `early-plugin-load` option, prior to `InnoDB` initialization. Failing to do so results in errors on startup and during `InnoDB` recovery.

To verify that the keyring plugin is active, use the `SHOW PLUGINS` statement or query the `INFORMATION_SCHEMA.PLUGINS` table. For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE 'keyring%';
+--------------+---------------+
| PLUGIN_NAME  | PLUGIN_STATUS |
+--------------+---------------+
| keyring_file | ACTIVE        |
+--------------+---------------+
```

- The `innodb_file_per_table` option must be enabled (the default). `InnoDB` tablespace encryption only supports file-per-table tablespaces. Alternatively, you can specify the `TABLESPACE='innodb_file_per_table'` option when creating an encrypted table or altering an existing table to enable encryption.
- Before using the `InnoDB` tablespace encryption feature, ensure that you have taken steps to prevent loss of the master encryption key. *If the master encryption key is lost, data stored in encrypted tablespace files is unrecoverable.* It is recommended that you create a backup of the `keyring` file immediately after creating the first encrypted table and before and after master key rotation. The `keyring` file location is defined by the `keyring_file_data` configuration option. For keyring plugin configuration information, see Section 6.5.4, “The MySQL Keyring”.

Enabling and Disabling InnoDB Tablespace Encryption

To enable encryption for a new `InnoDB` table, specify the `ENCRYPTION` option in a `CREATE TABLE` statement.

```
mysql> CREATE TABLE t1 (c1 INT) ENCRYPTION='Y';
```

To enable encryption for an existing `InnoDB` table, specify the `ENCRYPTION` option in an `ALTER TABLE` statement.

```
mysql> ALTER TABLE t1 ENCRYPTION='Y';
```

To disable encryption for an `InnoDB` table, set `ENCRYPTION='N'` using `ALTER TABLE`.

```
mysql> ALTER TABLE t1 ENCRYPTION='N';
```

Note

Plan appropriately when altering an existing table with the `ENCRYPTION` option. `ALTER TABLE ... ENCRYPTION` operations rebuild the table using `ALGORITHM=COPY`. `ALGORITM=INPLACE` is not supported.

Redo Log Data Encryption

Redo log data encryption is enabled using the `innodb_redo_log_encrypt` configuration option. Redo log encryption is disabled by default.

As with tablespace data, redo log data encryption occurs when redo log data is written to disk, and decryption occurs when redo log data is read from disk. Once redo log data is read into memory, it is in unencrypted form. Redo log data is encrypted and decrypted using the tablepace encryption key.

When `innodb_redo_log_encrypt` is enabled, unencrypted redo log pages that are present on disk remain unencrypted, and new redo log pages are written to disk in encrypted form. Likewise, when `innodb_redo_log_encrypt` is disabled, encrypted redo log pages that are present on disk remain encrypted, and new redo log pages are written to disk in unencrypted form.

Redo log encryption metadata, including the tablespace encryption key, is stored in the header of the first redo log file (`ib_logfile0`). If this file is removed, redo log encryption is disabled.

Once redo log encryption is enabled, a normal restart without the keyring plugin or without the encryption key is not possible, as `InnoDB` must be able to scan redo pages during startup, which is not

possible if redo log pages are encrypted. Without the keyring plugin or the encryption key, only a forced startup without the redo logs (`SRV_FORCE_NO_LOG_REDO`) is possible. See Section 15.20.2, "Forcing InnoDB Recovery".

Undo Log Data Encryption

Undo log data encryption is enabled using the `innodb_undo_log_encrypt` configuration option. Undo log encryption only applies undo logs that reside in separate undo tablespaces. See Section 15.7.7, "Storing InnoDB Undo Logs in Separate Tablespaces". Encryption is not supported for undo log data that resides in the system tablespace. Undo log data encryption is disabled by default.

As with tablespace data, undo log data encryption occurs when undo log data is written to disk, and decryption occurs when undo log data is read from disk. Once undo log data is read into memory, it is in unencrypted form. Undo log data is encrypted and decrypted using the tablepace encryption key.

When `innodb_undo_log_encrypt` is enabled, unencrypted undo log pages that are present on disk remain unencrypted, and new undo log pages are written to disk in encrypted form. Likewise, when `innodb_undo_log_encrypt` is disabled, encrypted undo log pages that are present on disk remain encrypted, and new undo log pages are written to disk in unencrypted form.

Undo log encryption metadata, including the tablespace encryption key, is stored in the header of the undo log file (`undo*N*.ibd`, where *N* is the space ID).

InnoDB Tablespace Encryption and Master Key Rotation

The master encryption key should be rotated periodically and whenever you suspect that the key may have been compromised.

Master key rotation is an atomic, instance-level operation. Each time the master encryption key is rotated, all tablespace keys in the MySQL instance are re-encrypted and saved back to their respective tablespace headers. As an atomic operation, re-encryption must succeed for all tablespace keys once a rotation operation is initiated. If master key rotation is interrupted by a server failure, `InnoDB` rolls the operation forward on server restart. For more information, see InnoDB Tablespace Encryption and Recovery.

Rotating the master encryption key only changes the master encryption key and re-encrypts tablespace keys. It does not decrypt or re-encrypt associated tablespace data.

Rotating the master encryption key requires the `ENCRYPTION_KEY_ADMIN` or `SUPER` privilege.

To rotate the master encryption key, run:

```
mysql> ALTER INSTANCE ROTATE INNODB MASTER KEY;
```

`ALTER INSTANCE ROTATE INNODB MASTER KEY` supports concurrent DML. However, it cannot be run concurrently with `CREATE TABLE ... ENCRYPTED` or `ALTER TABLE ... ENCRYPTED` operations, and locks are taken to prevent conflicts that could arise from concurrent execution of these statements. If one of the conflicting statements is running, it must complete before another can proceed.

InnoDB Tablespace Encryption and Recovery

If a server failure occurs during master key rotation, `InnoDB` continues the operation on server restart.

The keyring plugin must be loaded prior to storage engine initialization so that the information necessary to decrypt tablespace data pages can be retrieved from tablespace headers before `InnoDB` initialization and recovery activities access tablespace data. (See InnoDB Tablespace Encryption Prerequisites.)

When `InnoDB` initialization and recovery begin, the master key rotation operation resumes. Due to the server failure, some tablespaces keys may already be encrypted using the new master encryption key. `InnoDB` reads the encryption data from each tablespace header, and if the data indicates that the

tablespace key is encrypted using the old master encryption key, `InnoDB` retrieves the old key from the keyring and uses it to decrypt the tablepace key. `InnoDB` then re-encrypts the tablespace key using the new master encryption key and saves the re-encrypted tablespace key back to the tablespace header.

Exporting Encrypted Tables

When an encrypted table is exported, `InnoDB` generates a *transfer key* that is used to encrypt the tablespace key. The encrypted tablespace key and transfer key are stored in a *`tablespace_name`*`.cfp` file. This file together with the encrypted tablespace file is required to perform an import operation. On import, `InnoDB` uses the transfer key to decrypt the tablespace key in the *`tablespace_name`*`.cfp` file. For related information, see Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance”.

InnoDB Tablespace Encryption and Replication

- The `ALTER INSTANCE ROTATE INNODB MASTER KEY` statement is only supported in replication environments where the master and slaves run a version of MySQL that supports the tablespace encryption feature.
- Successful `ALTER INSTANCE ROTATE INNODB MASTER KEY` statements are written to the binary log for replication on slaves.
- If an `ALTER INSTANCE ROTATE INNODB MASTER KEY` statement fails, it is not logged to the binary log and is not replicated on slaves.
- Replication of an `ALTER INSTANCE ROTATE INNODB MASTER KEY` operation fails if the keyring plugin is installed on the master but not on the slave.
- If the `keyring_file` plugin is installed on both the master and a slave but the slave does not have a `keyring` file, the replicated `ALTER INSTANCE ROTATE INNODB MASTER KEY` statement creates the `keyring` file on the slave, assuming the `keyring` file data is not cached in memory. `ALTER INSTANCE ROTATE INNODB MASTER KEY` uses `keyring` file data that is cached in memory, if available.

Identifying Tables that Use InnoDB Tablespace Encryption

When the `ENCRYPTION` option is specified in a `CREATE TABLE` or `ALTER TABLE` statement, it is recorded in the `CREATE_OPTIONS` field of `INFORMATION_SCHEMA.TABLES`. This field may be queried to identify encrypted tables in a MySQL instance.

```
mysql> SELECT TABLE_SCHEMA, TABLE_NAME, CREATE_OPTIONS FROM INFORMATION_SCHEMA.TABLES
    -> WHERE CREATE_OPTIONS LIKE '%ENCRYPTION="Y"%';
+--------------+------------+----------------+
| TABLE_SCHEMA | TABLE_NAME | CREATE_OPTIONS |
+--------------+------------+----------------+
| test         | t1         | ENCRYPTION="Y" |
+--------------+------------+----------------+
```

InnoDB Tablespace Encryption Usage Notes

- If the server exits or is stopped during normal operation, it is recommended to restart the server using the same encryption settings that were configured previously.
- The first master encryption key is generated when the first new or existing table is encrypted.
- Master key rotation re-encrypts tablespaces keys but does not change the tablespace key itself. To change a tablespace key, you must disable and re-enable table encryption using `ALTER TABLE` *`tbl_name`* `ENCRYPTION`, which is an `ALGORITHM=COPY` operation that rebuilds the table.
- If a table is created with both the `COMPRESSION` and `ENCRYPTION` options, compression is performed before tablespace data is encrypted.

- If a `keyring` file is empty or missing, the first execution of `ALTER INSTANCE ROTATE INNODB MASTER KEY` creates a master encryption key.
- Uninstalling the `keyring_file` plugin does not remove an existing `keyring` file.
- It is recommended that you not place the `keyring` file under the same directory as tablespace data files. The location of the `keyring` file is specified by the `keyring_file_data` option.
- Modifying the `keyring_file_data` option at runtime or restarting the server with a new `keyring_file_data` setting can cause previously encrypted tables to become inaccessible, resulting in the loss of data.

InnoDB Tablespace Encryption Limitations

- Advanced Encryption Standard (AES) is the only supported encryption algorithm. `InnoDB` tablespace encryption uses Electronic Codebook (ECB) block encryption mode for tablespace key encryption and Cipher Block Chaining (CBC) block encryption mode for data encryption.
- Altering the `ENCRYPTION` attribute of a table is an `ALGORITHM=COPY` operation. `ALGORITHM=INPLACE` is not supported.
- `InnoDB` tablespace encryption only supports `InnoDB` tables that are stored in a file-per-table tablespaces. Encryption is not supported for tables stored in other `InnoDB` tablespace types including general tablespaces, the system tablespace, undo log tablespaces, and the temporary tablespace.
- You cannot move or copy an encrypted table from a file-per-table tablespace to an unsupported `InnoDB` tablespace type.
- By default, tablespace encryption only applies to data in the tablespace. Redo log and undo log data may be encrypted using the `innodb_redo_log_encrypt` and `innodb_undo_log_encrypt` options. See Redo Log Data Encryption, and Undo Log Data Encryption. Binary log data is not encrypted.

15.8 InnoDB Tables and Indexes

This section covers topics related to `InnoDB` tables and indexes.

15.8.1 InnoDB Tables

This section covers topics related to `InnoDB` tables.

15.8.1.1 Creating InnoDB Tables

To create an `InnoDB` table, use the `CREATE TABLE` statement.

```
CREATE TABLE t1 (a INT, b CHAR (20), PRIMARY KEY (a)) ENGINE=InnoDB;
```

You do not need to specify the `ENGINE=InnoDB` clause if `InnoDB` is defined as the default storage engine, which it is by default. To check the default storage engine, issue the following statement:

```
mysql> SELECT @@default_storage_engine;
+--------------------------+
| @@default_storage_engine |
+--------------------------+
| InnoDB                   |
+--------------------------+
```

You might still use `ENGINE=InnoDB` clause if you plan to use `mysqldump` or replication to replay the `CREATE TABLE` statement on a server where the default storage engine is not `InnoDB`.

An InnoDB table and its indexes can be created in the system tablespace, in a file-per-table tablespace, or in a general tablespace. When innodb_file_per_table is enabled, which is the default, an InnoDB table is implicitly created in an individual file-per-table tablespace. Conversely, when innodb_file_per_table is disabled, an InnoDB table is implicitly created in the InnoDB system tablespace. To create a table in a general tablespace, use CREATE TABLE ... TABLESPACE syntax. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

When you create a table in a file-per-table tablespace, MySQL creates an .ibd tablespace file in a database directory under the MySQL data directory, by default. A table created in the InnoDB system tablespace is created in an existing ibdata file, which resides in the MySQL data directory. A table created in a general tablespace is created in an existing general tablespace .ibd file. General tablespace files can be created inside or outside of the MySQL data directory. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

Internally, InnoDB adds an entry for each table to the data dictionary. The entry includes the database name. For example, if table t1 is created in the test database, the data dictionary entry for the database name is 'test/t1'. This means you can create a table of the same name (t1) in a different database, and the table names do not collide inside InnoDB.

InnoDB Tables and Row Formats

The default row format for InnoDB tables is defined by the innodb_default_row_format configuration option, which has a default value of DYNAMIC. Dynamic and Compressed row format allow you to take advantage of InnoDB features such as table compression and efficient off-page storage of long column values. To use these row formats, innodb_file_per_table must be enabled (the default).

```
SET GLOBAL innodb_file_per_table=1;
CREATE TABLE t3 (a INT, b CHAR (20), PRIMARY KEY (a)) ROW_FORMAT=DYNAMIC;
CREATE TABLE t4 (a INT, b CHAR (20), PRIMARY KEY (a)) ROW_FORMAT=COMPRESSED;
```

Alternatively, you can use CREATE TABLE ... TABLESPACE syntax to create an InnoDB table in a general tablespace. General tablespaces support all row formats. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

```
CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=DYNAMIC;
```

CREATE TABLE ... TABLESPACE syntax can also be used to create InnoDB tables with a Dynamic row format in the system tablespace, alongside tables with a Compact or Redundant row format.

```
CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE = innodb_system ROW_FORMAT=DYNAMIC;
```

For more information about InnoDB row formats, see Section 15.10, “InnoDB Row Storage and Row Formats”. For how to determine the row format of an InnoDB table and the physical characteristics of InnoDB row formats, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

InnoDB Tables and Primary Keys

Always define a primary key for an InnoDB table, specifying the column or columns that:

- Are referenced by the most important queries.
- Are never left blank.
- Never have duplicate values.
- Rarely if ever change value once inserted.

For example, in a table containing information about people, you would not create a primary key on (firstname, lastname) because more than one person can have the same name, some people have blank last names, and sometimes people change their names. With so many constraints, often there is not an obvious set of columns to use as a primary key, so you create a new column with a

numeric ID to serve as all or part of the primary key. You can declare an auto-increment column so that ascending values are filled in automatically as rows are inserted:

```
-- The value of ID can act like a pointer between related items in different tables.
CREATE TABLE t5 (id INT AUTO_INCREMENT, b CHAR (20), PRIMARY KEY (id));

-- The primary key can consist of more than one column. Any autoinc column must come first.
CREATE TABLE t6 (id INT AUTO_INCREMENT, a INT, b CHAR (20), PRIMARY KEY (id,a));
```

Although the table works correctly without defining a primary key, the primary key is involved with many aspects of performance and is a crucial design aspect for any large or frequently used table. It is recommended that you always specify a primary key in the `CREATE TABLE` statement. If you create the table, load data, and then run `ALTER TABLE` to add a primary key later, that operation is much slower than defining the primary key when creating the table.

Viewing InnoDB Table Properties

To view the properties of an `InnoDB` table, issue a `SHOW TABLE STATUS` statement:

```
mysql> SHOW TABLE STATUS FROM test LIKE 't%' \G;
*************************** 1. row ***************************
           Name: t1
         Engine: InnoDB
        Version: 10
     Row_format: Compact
           Rows: 0
 Avg_row_length: 0
    Data_length: 16384
Max_data_length: 0
   Index_length: 0
      Data_free: 0
 Auto_increment: NULL
    Create_time: 2015-03-16 15:13:31
    Update_time: NULL
     Check_time: NULL
      Collation: utf8mb4_0900_ai_ci
       Checksum: NULL
 Create_options:
        Comment:
```

For information about `SHOW TABLE STATUS` output, see Section 13.7.5.36, “SHOW TABLE STATUS Syntax”.

`InnoDB` table properties may also be queried using the `InnoDB` Information Schema system tables:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME='test/t1' \G
*************************** 1. row ***************************
     TABLE_ID: 45
         NAME: test/t1
         FLAG: 1
       N_COLS: 5
        SPACE: 35
   ROW_FORMAT: Compact
ZIP_PAGE_SIZE: 0
   SPACE_TYPE: Single
```

For more information, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

15.8.1.2 The Physical Row Structure of an InnoDB Table

The physical row structure of an `InnoDB` table depends on the row format specified when the table is created. If a row format is not specified, the default row format is used. The default row format for `InnoDB` tables is defined by the `innodb_default_row_format` configuration option, which has a default value of `DYNAMIC`.

The following sections describe the characteristics of `InnoDB` row formats.

- Determining the Row Format of an InnoDB Table
- Redundant Row Format Characteristics
- COMPACT Row Format Characteristics
- DYNAMIC and COMPRESSED Row Format Characteristics

For more information about InnoDB row formats, see Section 15.10, “InnoDB Row Storage and Row Formats”.

Determining the Row Format of an InnoDB Table

To determine the row format of an InnoDB table, you can use SHOW TABLE STATUS. For example:

```
mysql> SHOW TABLE STATUS IN test1\G
*************************** 1. row ***************************
           Name: t1
         Engine: InnoDB
        Version: 10
     Row_format: Dynamic
           Rows: 0
 Avg_row_length: 0
    Data_length: 16384
Max_data_length: 0
   Index_length: 16384
      Data_free: 0
 Auto_increment: 1
    Create_time: 2016-09-14 16:29:38
    Update_time: NULL
     Check_time: NULL
      Collation: utf8mb4_0900_ai_ci
       Checksum: NULL
 Create_options:
        Comment:
```

You can also determine the row format of an InnoDB table by querying INFORMATION_SCHEMA.INNODB_SYS_TABLES.

```
mysql> SELECT NAME, ROW_FORMAT FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME='test1/t1';
+----------+------------+
| NAME     | ROW_FORMAT |
+----------+------------+
| test1/t1 | Dynamic    |
+----------+------------+
```

Redundant Row Format Characteristics

The REDUNDANT format is available to retain compatibility with older versions of MySQL.

Rows in InnoDB tables that use REDUNDANT row format have the following characteristics:

- Each index record contains a 6-byte header. The header is used to link together consecutive records, and also in row-level locking.
- Records in the clustered index contain fields for all user-defined columns. In addition, there is a 6-byte transaction ID field and a 7-byte roll pointer field.
- If no primary key was defined for a table, each clustered index record also contains a 6-byte row ID field.
- Each secondary index record also contains all the primary key fields defined for the clustered index key that are not in the secondary index.
- A record contains a pointer to each field of the record. If the total length of the fields in a record is less than 128 bytes, the pointer is one byte; otherwise, two bytes. The array of these pointers is called the record directory. The area where these pointers point is called the data part of the record.

- Internally, `InnoDB` stores fixed-length character columns such as `CHAR(10)` in a fixed-length format. `InnoDB` does not truncate trailing spaces from `VARCHAR` columns.
- `InnoDB` encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored off-page. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.
- An SQL `NULL` value reserves one or two bytes in the record directory. Besides that, an SQL `NULL` value reserves zero bytes in the data part of the record if stored in a variable length column. In a fixed-length column, it reserves the fixed length of the column in the data part of the record. Reserving the fixed space for `NULL` values enables an update of the column from `NULL` to a non-`NULL` value to be done in place without causing fragmentation of the index page.

COMPACT Row Format Characteristics

The `COMPACT` row format decreases row storage space by about 20% compared to the `REDUNDANT` format at the cost of increasing CPU use for some operations. If your workload is a typical one that is limited by cache hit rates and disk speed, `COMPACT` format is likely to be faster. If the workload is a rare case that is limited by CPU speed, compact format might be slower.

Rows in `InnoDB` tables that use `COMPACT` row format have the following characteristics:

- Each index record contains a 5-byte header that may be preceded by a variable-length header. The header is used to link together consecutive records, and also in row-level locking.
- The variable-length part of the record header contains a bit vector for indicating `NULL` columns. If the number of columns in the index that can be `NULL` is *`N`*, the bit vector occupies `CEILING(`*`N`*`/8)` bytes. (For example, if there are anywhere from 9 to 15 columns that can be `NULL`, the bit vector uses two bytes.) Columns that are `NULL` do not occupy space other than the bit in this vector. The variable-length part of the header also contains the lengths of variable-length columns. Each length takes one or two bytes, depending on the maximum length of the column. If all columns in the index are `NOT NULL` and have a fixed length, the record header has no variable-length part.
- For each non-`NULL` variable-length field, the record header contains the length of the column in one or two bytes. Two bytes are only needed if part of the column is stored externally in overflow pages or the maximum length exceeds 255 bytes and the actual length exceeds 127 bytes. For an externally stored column, the 2-byte length indicates the length of the internally stored part plus the 20-byte pointer to the externally stored part. The internal part is 768 bytes, so the length is 768+20. The 20-byte pointer stores the true length of the column.
- The record header is followed by the data contents of the non-`NULL` columns.
- Records in the clustered index contain fields for all user-defined columns. In addition, there is a 6-byte transaction ID field and a 7-byte roll pointer field.
- If no primary key was defined for a table, each clustered index record also contains a 6-byte row ID field.
- Each secondary index record also contains all the primary key fields defined for the clustered index key that are not in the secondary index. If any of these primary key fields are variable length, the record header for each secondary index has a variable-length part to record their lengths, even if the secondary index is defined on fixed-length columns.
- Internally, for nonvariable-length character sets, `InnoDB` stores fixed-length character columns such as `CHAR(10)` in a fixed-length format.

  `InnoDB` does not truncate trailing spaces from `VARCHAR` columns.
- Internally, for variable-length character sets such as `utf8mb3` and `utf8mb4`, `InnoDB` attempts to store `CHAR(`*`N`*`)` in *`N`* bytes by trimming trailing spaces. If the byte length of a `CHAR(`*`N`*`)` column value exceeds *`N`* bytes, `InnoDB` trims trailing spaces to a minimum of the column value byte length. The maximum length of a `CHAR(`*`N`*`)` column is the maximum character byte length × *`N`*.

`InnoDB` reserves a minimum of *`N`* bytes for `CHAR(N)`. Reserving the minimum space *`N`* in many cases enables column updates to be done in place without causing fragmentation of the index page. By comparison, for `ROW_FORMAT=REDUNDANT`, `CHAR(N)` columns occupy the maximum character byte length × *`N`*.

`InnoDB` encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored off-page. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.

`ROW_FORMAT=DYNAMIC` and `ROW_FORMAT=COMPRESSED` handle `CHAR` storage in the same way as `ROW_FORMAT=COMPACT`.

DYNAMIC and COMPRESSED Row Format Characteristics

`DYNAMIC` and `COMPRESSED` row formats are variations of the `COMPACT` row format. For information about these row formats, see Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats”.

15.8.1.3 Moving or Copying InnoDB Tables

This section describes techniques for moving or copying some or all `InnoDB` tables to a different server or instance. For example, you might move an entire MySQL instance to a larger, faster server; you might clone an entire MySQL instance to a new replication slave server; you might copy individual tables to another instance to develop and test an application, or to a data warehouse server to produce reports.

On Windows, `InnoDB` always stores database and table names internally in lowercase. To move databases in a binary format from Unix to Windows or from Windows to Unix, create all databases and tables using lowercase names. A convenient way to accomplish this is to add the following line to the `[mysqld]` section of your `my.cnf` or `my.ini` file before creating any databases or tables:

```
[mysqld]
lower_case_table_names=1
```

Techniques for moving or copying `InnoDB` tables include:

- Transportable Tablespaces
- MySQL Enterprise Backup
- Copying Data Files (Cold Backup Method)
- Export and Import (mysqldump)

Transportable Tablespaces

The transportable tablespaces feature uses `FLUSH TABLES ... FOR EXPORT` to ready `InnoDB` tables for copying from one server instance to another. To use this feature, `InnoDB` tables must be created with `innodb_file_per_table` set to `ON` so that each `InnoDB` table has its own tablespace. For usage information, see Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance”.

MySQL Enterprise Backup

The MySQL Enterprise Backup product lets you back up a running MySQL database, including `InnoDB` and `MyISAM` tables, with minimal disruption to operations while producing a consistent snapshot of the database. When MySQL Enterprise Backup is copying `InnoDB` tables, reads and writes to both `InnoDB` and `MyISAM` tables can continue. During the copying of `MyISAM` and other non-InnoDB tables, reads (but not writes) to those tables are permitted. In addition, MySQL Enterprise Backup can create compressed backup files, and back up subsets of `InnoDB` tables. In conjunction with the MySQL binary log, you can perform point-in-time recovery. MySQL Enterprise Backup is included as part of the MySQL Enterprise subscription.

For more details about MySQL Enterprise Backup, see Section 29.2, “MySQL Enterprise Backup Overview”.

Copying Data Files (Cold Backup Method)

You can move an `InnoDB` database simply by copying all the relevant files listed under "Cold Backups" in Section 15.17.1, “InnoDB Backup”.

Like `MyISAM` data files, `InnoDB` data and log files are binary-compatible on all platforms having the same floating-point number format. If the floating-point formats differ but you have not used `FLOAT` or `DOUBLE` data types in your tables, then the procedure is the same: simply copy the relevant files.

When you move or copy file-per-table `.ibd` files, the database directory name must be the same on the source and destination systems. The table definition stored in the `InnoDB` shared tablespace includes the database name. The transaction IDs and log sequence numbers stored in the tablespace files also differ between databases.

To move an `.ibd` file and the associated table from one database to another, use a `RENAME TABLE` statement:

```
RENAME TABLE db1.tbl_name TO db2.tbl_name;
```

If you have a “clean” backup of an `.ibd` file, you can restore it to the MySQL installation from which it originated as follows:

1. The table must not have been dropped or truncated since you copied the `.ibd` file, because doing so changes the table ID stored inside the tablespace.
2. Issue this `ALTER TABLE` statement to delete the current `.ibd` file:

   ```
   ALTER TABLE tbl_name DISCARD TABLESPACE;
   ```
3. Copy the backup `.ibd` file to the proper database directory.
4. Issue this `ALTER TABLE` statement to tell `InnoDB` to use the new `.ibd` file for the table:

   ```
   ALTER TABLE tbl_name IMPORT TABLESPACE;
   ```

Note

The `ALTER TABLE ... IMPORT TABLESPACE` feature does not enforce foreign key constraints on imported data.

In this context, a “clean” `.ibd` file backup is one for which the following requirements are satisfied:

- There are no uncommitted modifications by transactions in the `.ibd` file.
- There are no unmerged insert buffer entries in the `.ibd` file.
- Purge has removed all delete-marked index records from the `.ibd` file.
- `mysqld` has flushed all modified pages of the `.ibd` file from the buffer pool to the file.

You can make a clean backup `.ibd` file using the following method:

1. Stop all activity from the `mysqld` server and commit all transactions.
2. Wait until `SHOW ENGINE INNODB STATUS` shows that there are no active transactions in the database, and the main thread status of `InnoDB` is `Waiting for server activity`. Then you can make a copy of the `.ibd` file.

Another method for making a clean copy of an `.ibd` file is to use the MySQL Enterprise Backup product:

1. Use MySQL Enterprise Backup to back up the `InnoDB` installation.
2. Start a second `mysqld` server on the backup and let it clean up the `.ibd` files in the backup.

Export and Import (mysqldump)

You can use `mysqldump` to dump your tables on one machine and then import the dump files on the other machine. Using this method, it does not matter whether the formats differ or if your tables contain floating-point data.

One way to increase the performance of this method is to switch off autocommit mode when importing data, assuming that the tablespace has enough space for the big rollback segment that the import transactions generate. Do the commit only after importing a whole table or a segment of a table.

15.8.1.4 Converting Tables from MyISAM to InnoDB

If you have `MyISAM` tables that you want to convert to `InnoDB` for better reliability and scalability, review the following guidelines and tips before converting.

Note

Partitioned `MyISAM` tables created in previous versions of MySQL are not compatible with MySQL 8.0. Such tables must be prepared prior to upgrade, either by removing the partitioning, or by converting them to `InnoDB`. See Section 22.6.2, “Partitioning Limitations Relating to Storage Engines”, for more information.

- Adjusting Memory Usage for MyISAM and InnoDB
- Handling Too-Long Or Too-Short Transactions
- Handling Deadlocks
- Planning the Storage Layout
- Converting an Existing Table
- Cloning the Structure of a Table
- Transferring Existing Data
- Storage Requirements
- Defining a PRIMARY KEY for Each Table
- Application Performance Considerations
- Understanding Files Associated with InnoDB Tables

Adjusting Memory Usage for MyISAM and InnoDB

As you transition away from `MyISAM` tables, lower the value of the `key_buffer_size` configuration option to free memory no longer needed for caching results. Increase the value of the `innodb_buffer_pool_size` configuration option, which performs a similar role of allocating cache memory for `InnoDB` tables. The `InnoDB` buffer pool caches both table data and index data, speeding up lookups for queries and keeping query results in memory for reuse. For guidance regarding buffer pool size configuration, see Section 8.12.3.1, “How MySQL Uses Memory”.

On a busy server, run benchmarks with the query cache turned off. The `InnoDB` buffer pool provides similar benefits, so the query cache might be tying up memory unnecessarily. For information about the query cache, see Section 8.10.3, “The MySQL Query Cache”.

Handling Too-Long Or Too-Short Transactions

Because MyISAM tables do not support transactions, you might not have paid much attention to the autocommit configuration option and the COMMIT and ROLLBACK statements. These keywords are important to allow multiple sessions to read and write InnoDB tables concurrently, providing substantial scalability benefits in write-heavy workloads.

While a transaction is open, the system keeps a snapshot of the data as seen at the beginning of the transaction, which can cause substantial overhead if the system inserts, updates, and deletes millions of rows while a stray transaction keeps running. Thus, take care to avoid transactions that run for too long:

- If you are using a mysql session for interactive experiments, always COMMIT (to finalize the changes) or ROLLBACK (to undo the changes) when finished. Close down interactive sessions rather than leave them open for long periods, to avoid keeping transactions open for long periods by accident.
- Make sure that any error handlers in your application also ROLLBACK incomplete changes or COMMIT completed changes.
- ROLLBACK is a relatively expensive operation, because INSERT, UPDATE, and DELETE operations are written to InnoDB tables prior to the COMMIT, with the expectation that most changes are committed successfully and rollbacks are rare. When experimenting with large volumes of data, avoid making changes to large numbers of rows and then rolling back those changes.
- When loading large volumes of data with a sequence of INSERT statements, periodically COMMIT the results to avoid having transactions that last for hours. In typical load operations for data warehousing, if something goes wrong, you truncate the table (using TRUNCATE TABLE) and start over from the beginning rather than doing a ROLLBACK.

The preceding tips save memory and disk space that can be wasted during too-long transactions. When transactions are shorter than they should be, the problem is excessive I/O. With each COMMIT, MySQL makes sure each change is safely recorded to disk, which involves some I/O.

- For most operations on InnoDB tables, you should use the setting autocommit=0. From an efficiency perspective, this avoids unnecessary I/O when you issue large numbers of consecutive INSERT, UPDATE, or DELETE statements. From a safety perspective, this allows you to issue a ROLLBACK statement to recover lost or garbled data if you make a mistake on the mysql command line, or in an exception handler in your application.
- The time when autocommit=1 is suitable for InnoDB tables is when running a sequence of queries for generating reports or analyzing statistics. In this situation, there is no I/O penalty related to COMMIT or ROLLBACK, and InnoDB can automatically optimize the read-only workload.
- If you make a series of related changes, finalize all the changes at once with a single COMMIT at the end. For example, if you insert related pieces of information into several tables, do a single COMMIT after making all the changes. Or if you run many consecutive INSERT statements, do a single COMMIT after all the data is loaded; if you are doing millions of INSERT statements, perhaps split up the huge transaction by issuing a COMMIT every ten thousand or hundred thousand records, so the transaction does not grow too large.
- Remember that even a SELECT statement opens a transaction, so after running some report or debugging queries in an interactive mysql session, either issue a COMMIT or close the mysql session.

Handling Deadlocks

You might see warning messages referring to “deadlocks” in the MySQL error log, or the output of SHOW ENGINE INNODB STATUS. Despite the scary-sounding name, a deadlock is not a serious issue for InnoDB tables, and often does not require any corrective action. When two transactions start modifying multiple tables, accessing the tables in a different order, they can reach a state where each

transaction is waiting for the other and neither can proceed. When deadlock detection is enabled (the default), MySQL immediately detects this condition and cancels (rolls back) the “smaller” transaction, allowing the other to proceed. If deadlock detection is disabled using the `innodb_deadlock_detect` configuration option, `InnoDB` relies on the `innodb_lock_wait_timeout` setting to roll back transactions in case of a deadlock.

Either way, your applications need error-handling logic to restart a transaction that is forcibly cancelled due to a deadlock. When you re-issue the same SQL statements as before, the original timing issue no longer applies. Either the other transaction has already finished and yours can proceed, or the other transaction is still in progress and your transaction waits until it finishes.

If deadlock warnings occur constantly, you might review the application code to reorder the SQL operations in a consistent way, or to shorten the transactions. You can test with the `innodb_print_all_deadlocks` option enabled to see all deadlock warnings in the MySQL error log, rather than only the last warning in the `SHOW ENGINE INNODB STATUS` output.

For more information, see Section 15.5.5, “Deadlocks in InnoDB”.

Planning the Storage Layout

To get the best performance from `InnoDB` tables, you can adjust a number of parameters related to storage layout.

When you convert `MyISAM` tables that are large, frequently accessed, and hold vital data, investigate and consider the `innodb_file_per_table` and `innodb_page_size` configuration options, and the `ROW_FORMAT` and `KEY_BLOCK_SIZE` clauses of the `CREATE TABLE` statement.

During your initial experiments, the most important setting is `innodb_file_per_table`. When this setting is enabled, which is the default, new `InnoDB` tables are implicitly created in file-per-table tablespaces. In contrast with the `InnoDB` system tablespace, file-per-table tablespaces allow disk space to be reclaimed by the operating system when a table is truncated or dropped. File-per-table tablespaces also support DYNAMIC and COMPRESSED row formats and associated features such as table compression, efficient off-page storage for long variable-length columns, and large index prefixes. For more information, see Section 15.7.4, “InnoDB File-Per-Table Tablespaces”.

You can also store `InnoDB` tables in a shared general tablespace, which support multiple tables and all row formats. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

Converting an Existing Table

To convert a non-`InnoDB` table to use `InnoDB` use `ALTER TABLE`:

```
ALTER TABLE table_name ENGINE=InnoDB;
```

Cloning the Structure of a Table

You might make an `InnoDB` table that is a clone of a MyISAM table, rather than using `ALTER TABLE` to perform conversion, to test the old and new table side-by-side before switching.

Create an empty `InnoDB` table with identical column and index definitions. Use `SHOW CREATE TABLE table_name\G` to see the full `CREATE TABLE` statement to use. Change the `ENGINE` clause to `ENGINE=INNODB`.

Transferring Existing Data

To transfer a large volume of data into an empty `InnoDB` table created as shown in the previous section, insert the rows with `INSERT INTO innodb_table SELECT * FROM myisam_table ORDER BY primary_key_columns`.

You can also create the indexes for the `InnoDB` table after inserting the data. Historically, creating new secondary indexes was a slow operation for InnoDB, but now you can create the indexes after the data is loaded with relatively little overhead from the index creation step.

If you have UNIQUE constraints on secondary keys, you can speed up a table import by turning off the uniqueness checks temporarily during the import operation:

```
SET unique_checks=0;
... import operation ...
SET unique_checks=1;
```

For big tables, this saves disk I/O because InnoDB can use its change buffer to write secondary index records as a batch. Be certain that the data contains no duplicate keys. unique_checks permits but does not require storage engines to ignore duplicate keys.

For better control over the insertion process, you can insert big tables in pieces:

```
INSERT INTO newtable SELECT * FROM oldtable
   WHERE yourkey > something AND yourkey <= somethingelse;
```

After all records are inserted, you can rename the tables.

During the conversion of big tables, increase the size of the InnoDB buffer pool to reduce disk I/O, to a maximum of 80% of physical memory. You can also increase the size of InnoDB log files.

Storage Requirements

If you intend to make several temporary copies of your data in InnoDB tables during the conversion process, it is recommended that you create the tables in file-per-table tablespaces so that you can reclaim the disk space when you drop the tables. When the innodb_file_per_table configuration option is enabled (the default), newly created InnoDB tables are implicitly created in file-per-table tablespaces.

Whether you convert the MyISAM table directly or create a cloned InnoDB table, make sure that you have sufficient disk space to hold both the old and new tables during the process. **InnoDB tables require more disk space than MyISAM tables.** If an ALTER TABLE operation runs out of space, it starts a rollback, and that can take hours if it is disk-bound. For inserts, InnoDB uses the insert buffer to merge secondary index records to indexes in batches. That saves a lot of disk I/O. For rollback, no such mechanism is used, and the rollback can take 30 times longer than the insertion.

In the case of a runaway rollback, if you do not have valuable data in your database, it may be advisable to kill the database process rather than wait for millions of disk I/O operations to complete. For the complete procedure, see Section 15.20.2, “Forcing InnoDB Recovery”.

Defining a PRIMARY KEY for Each Table

The PRIMARY KEY clause is a critical factor affecting the performance of MySQL queries and the space usage for tables and indexes. The primary key uniquely identifies a row in a table. Every row in the table must have a primary key value, and no two rows can have the same primary key value.

These are guidelines for the primary key, followed by more detailed explanations.

- Declare a PRIMARY KEY for each table. Typically, it is the most important column that you refer to in WHERE clauses when looking up a single row.
- Declare the PRIMARY KEY clause in the original CREATE TABLE statement, rather than adding it later through an ALTER TABLE statement.
- Choose the column and its data type carefully. Prefer numeric columns over character or string ones.
- Consider using an auto-increment column if there is not another stable, unique, non-null, numeric column to use.
- An auto-increment column is also a good choice if there is any doubt whether the value of the primary key column could ever change. Changing the value of a primary key column is an expensive operation, possibly involving rearranging data within the table and within each secondary index.

Consider adding a primary key to any table that does not already have one. Use the smallest practical numeric type based on the maximum projected size of the table. This can make each row slightly more compact, which can yield substantial space savings for large tables. The space savings are multiplied if the table has any secondary indexes, because the primary key value is repeated in each secondary index entry. In addition to reducing data size on disk, a small primary key also lets more data fit into the buffer pool, speeding up all kinds of operations and improving concurrency.

If the table already has a primary key on some longer column, such as a `VARCHAR`, consider adding a new unsigned `AUTO_INCREMENT` column and switching the primary key to that, even if that column is not referenced in queries. This design change can produce substantial space savings in the secondary indexes. You can designate the former primary key columns as `UNIQUE NOT NULL` to enforce the same constraints as the `PRIMARY KEY` clause, that is, to prevent duplicate or null values across all those columns.

If you spread related information across multiple tables, typically each table uses the same column for its primary key. For example, a personnel database might have several tables, each with a primary key of employee number. A sales database might have some tables with a primary key of customer number, and other tables with a primary key of order number. Because lookups using the primary key are very fast, you can construct efficient join queries for such tables.

If you leave the `PRIMARY KEY` clause out entirely, MySQL creates an invisible one for you. It is a 6-byte value that might be longer than you need, thus wasting space. Because it is hidden, you cannot refer to it in queries.

Application Performance Considerations

The reliability and scalability features of `InnoDB` require more disk storage than equivalent `MyISAM` tables. You might change the column and index definitions slightly, for better space utilization, reduced I/O and memory consumption when processing result sets, and better query optimization plans making efficient use of index lookups.

If you do set up a numeric ID column for the primary key, use that value to cross-reference with related values in any other tables, particularly for join queries. For example, rather than accepting a country name as input and doing queries searching for the same name, do one lookup to determine the country ID, then do other queries (or a single join query) to look up relevant information across several tables. Rather than storing a customer or catalog item number as a string of digits, potentially using up several bytes, convert it to a numeric ID for storing and querying. A 4-byte unsigned `INT` column can index over 4 billion items (with the US meaning of billion: 1000 million). For the ranges of the different integer types, see Section 11.2.1, “Integer Types (Exact Value) - INTEGER, INT, SMALLINT, TINYINT, MEDIUMINT, BIGINT”.

Understanding Files Associated with InnoDB Tables

`InnoDB` files require more care and planning than `MyISAM` files do.

- You must not delete the ibdata files that represent the `InnoDB` system tablespace.
- Methods of moving or copying `InnoDB` tables to a different server are described in Section 15.8.1.3, “Moving or Copying InnoDB Tables”.

15.8.1.5 AUTO_INCREMENT Handling in InnoDB

`InnoDB` provides a configurable locking mechanism that can significantly improve scalability and performance of SQL statements that add rows to tables with `AUTO_INCREMENT` columns. To use the `AUTO_INCREMENT` mechanism with an `InnoDB` table, an `AUTO_INCREMENT` column must be defined as part of an index such that it is possible to perform the equivalent of an indexed `SELECT MAX(ai_col)` lookup on the table to obtain the maximum column value. Typically, this is achieved by making the column the first column of some table index.

This section describes the behavior of `AUTO_INCREMENT` lock modes, usage implications for different `AUTO_INCREMENT` lock mode settings, and how `InnoDB` initializes the `AUTO_INCREMENT` counter.

- InnoDB AUTO_INCREMENT Lock Modes
- InnoDB AUTO_INCREMENT Lock Mode Usage Implications
- InnoDB AUTO_INCREMENT Counter Initialization

InnoDB AUTO_INCREMENT Lock Modes

This section describes the behavior of `AUTO_INCREMENT` lock modes used to generate auto-increment values, and how each lock mode affects replication. Auto-increment lock modes are configured at startup using the `innodb_autoinc_lock_mode` configuration parameter.

The following terms are used in describing `innodb_autoinc_lock_mode` settings:

- "`INSERT`-like" statements

  All statements that generate new rows in a table, including `INSERT`, `INSERT ... SELECT`, `REPLACE`, `REPLACE ... SELECT`, and `LOAD DATA`. Includes "simple-inserts", "bulk-inserts", and "mixed-mode" inserts.

- "Simple inserts"

  Statements for which the number of rows to be inserted can be determined in advance (when the statement is initially processed). This includes single-row and multiple-row `INSERT` and `REPLACE` statements that do not have a nested subquery, but not `INSERT ... ON DUPLICATE KEY UPDATE`.

- "Bulk inserts"

  Statements for which the number of rows to be inserted (and the number of required auto-increment values) is not known in advance. This includes `INSERT ... SELECT`, `REPLACE ... SELECT`, and `LOAD DATA` statements, but not plain `INSERT`. `InnoDB` assigns new values for the `AUTO_INCREMENT` column one at a time as each row is processed.

- "Mixed-mode inserts"

  These are "simple insert" statements that specify the auto-increment value for some (but not all) of the new rows. An example follows, where `c1` is an `AUTO_INCREMENT` column of table `t1`:

  ```
  INSERT INTO t1 (c1,c2) VALUES (1,'a'), (NULL,'b'), (5,'c'), (NULL,'d');
  ```

  Another type of "mixed-mode insert" is `INSERT ... ON DUPLICATE KEY UPDATE`, which in the worst case is in effect an `INSERT` followed by a `UPDATE`, where the allocated value for the `AUTO_INCREMENT` column may or may not be used during the update phase.

There are three possible settings for the `innodb_autoinc_lock_mode` configuration parameter. The settings are 0, 1, or 2, for "traditional", "consecutive", or "interleaved" lock mode, respectively.

- `innodb_autoinc_lock_mode = 0` ("traditional" lock mode)

  The traditional lock mode provides the same behavior that existed before the `innodb_autoinc_lock_mode` configuration parameter was introduced in MySQL 5.1. The traditional lock mode option is provided for backward compatibility, performance testing, and working around issues with "mixed-mode inserts", due to possible differences in semantics.

  In this lock mode, all "INSERT-like" statements obtain a special table-level `AUTO-INC` lock for inserts into tables with `AUTO_INCREMENT` columns. This lock is normally held to the end of the statement (not to the end of the transaction) to ensure that auto-increment values are assigned in a predictable and repeatable order for a given sequence of `INSERT` statements, and to ensure that auto-increment values assigned by any given statement are consecutive.

In the case of statement-based replication, this means that when an SQL statement is replicated on a slave server, the same values are used for the auto-increment column as on the master server. The result of execution of multiple `INSERT` statements is deterministic, and the slave reproduces the same data as on the master. If auto-increment values generated by multiple `INSERT` statements were interleaved, the result of two concurrent `INSERT` statements would be nondeterministic, and could not reliably be propagated to a slave server using statement-based replication.

To make this clear, consider an example that uses this table:

```
CREATE TABLE t1 (
  c1 INT(11) NOT NULL AUTO_INCREMENT,
  c2 VARCHAR(10) DEFAULT NULL,
  PRIMARY KEY (c1)
) ENGINE=InnoDB;
```

Suppose that there are two transactions running, each inserting rows into a table with an `AUTO_INCREMENT` column. One transaction is using an `INSERT ... SELECT` statement that inserts 1000 rows, and another is using a simple `INSERT` statement that inserts one row:

```
Tx1: INSERT INTO t1 (c2) SELECT 1000 rows from another table ...
Tx2: INSERT INTO t1 (c2) VALUES ('xxx');
```

`InnoDB` cannot tell in advance how many rows are retrieved from the `SELECT` in the `INSERT` statement in Tx1, and it assigns the auto-increment values one at a time as the statement proceeds. With a table-level lock, held to the end of the statement, only one `INSERT` statement referring to table `t1` can execute at a time, and the generation of auto-increment numbers by different statements is not interleaved. The auto-increment value generated by the Tx1 `INSERT ... SELECT` statement are consecutive, and the (single) auto-increment value used by the `INSERT` statement in Tx2 are either smaller or larger than all those used for Tx1, depending on which statement executes first.

As long as the SQL statements execute in the same order when replayed from the binary log (when using statement-based replication, or in recovery scenarios), the results are the same as they were when Tx1 and Tx2 first ran. Thus, table-level locks held until the end of a statement make `INSERT` statements using auto-increment safe for use with statement-based replication. However, those table-level locks limit concurrency and scalability when multiple transactions are executing insert statements at the same time.

In the preceding example, if there were no table-level lock, the value of the auto-increment column used for the `INSERT` in Tx2 depends on precisely when the statement executes. If the `INSERT` of Tx2 executes while the `INSERT` of Tx1 is running (rather than before it starts or after it completes), the specific auto-increment values assigned by the two `INSERT` statements are nondeterministic, and may vary from run to run.

Under the consecutive lock mode, `InnoDB` can avoid using table-level `AUTO-INC` locks for “simple insert” statements where the number of rows is known in advance, and still preserve deterministic execution and safety for statement-based replication.

If you are not using the binary log to replay SQL statements as part of recovery or replication, the interleaved lock mode can be used to eliminate all use of table-level `AUTO-INC` locks for even greater concurrency and performance, at the cost of permitting gaps in auto-increment numbers assigned by a statement and potentially having the numbers assigned by concurrently executing statements interleaved.

- `innodb_autoinc_lock_mode = 1` (“consecutive” lock mode)

  This is the default lock mode. In this mode, “bulk inserts” use the special `AUTO-INC` table-level lock and hold it until the end of the statement. This applies to all `INSERT ... SELECT`, `REPLACE ... SELECT`, and `LOAD DATA` statements. Only one statement holding the `AUTO-INC` lock can execute

at a time. If the source table of the bulk insert operation is different from the target table, the AUTO-INC lock on the target table is taken after a shared lock is taken on the first row selected from the source table. If the source and target of the bulk insert operation are the same table, the AUTO-INC lock is taken after shared locks are taken on all selected rows.

"Simple inserts" (for which the number of rows to be inserted is known in advance) avoid table-level AUTO-INC locks by obtaining the required number of auto-increment values under the control of a mutex (a light-weight lock) that is only held for the duration of the allocation process, *not* until the statement completes. No table-level AUTO-INC lock is used unless an AUTO-INC lock is held by another transaction. If another transaction holds an AUTO-INC lock, a "simple insert" waits for the AUTO-INC lock, as if it were a "bulk insert".

This lock mode ensures that, in the presence of INSERT statements where the number of rows is not known in advance (and where auto-increment numbers are assigned as the statement progresses), all auto-increment values assigned by any "INSERT-like" statement are consecutive, and operations are safe for statement-based replication.

Simply put, this lock mode significantly improves scalability while being safe for use with statement-based replication. Further, as with "traditional" lock mode, auto-increment numbers assigned by any given statement are *consecutive*. There is *no change* in semantics compared to "traditional" mode for any statement that uses auto-increment, with one important exception.

The exception is for "mixed-mode inserts", where the user provides explicit values for an AUTO_INCREMENT column for some, but not all, rows in a multiple-row "simple insert". For such inserts, InnoDB allocates more auto-increment values than the number of rows to be inserted. However, all values automatically assigned are consecutively generated (and thus higher than) the auto-increment value generated by the most recently executed previous statement. "Excess" numbers are lost.

- `innodb_autoinc_lock_mode = 2` ("interleaved" lock mode)

  In this lock mode, no "INSERT-like" statements use the table-level AUTO-INC lock, and multiple statements can execute at the same time. This is the fastest and most scalable lock mode, but it is *not safe* when using statement-based replication or recovery scenarios when SQL statements are replayed from the binary log.

  In this lock mode, auto-increment values are guaranteed to be unique and monotonically increasing across all concurrently executing "INSERT-like" statements. However, because multiple statements can be generating numbers at the same time (that is, allocation of numbers is *interleaved* across statements), the values generated for the rows inserted by any given statement may not be consecutive.

  If the only statements executing are "simple inserts" where the number of rows to be inserted is known ahead of time, there are no gaps in the numbers generated for a single statement, except for "mixed-mode inserts". However, when "bulk inserts" are executed, there may be gaps in the auto-increment values assigned by any given statement.

InnoDB AUTO_INCREMENT Lock Mode Usage Implications

- Using auto-increment with replication

  If you are using statement-based replication, set innodb_autoinc_lock_mode to 0 or 1 and use the same value on the master and its slaves. Auto-increment values are not ensured to be the same on the slaves as on the master if you use innodb_autoinc_lock_mode = 2 ("interleaved") or configurations where the master and slaves do not use the same lock mode.

  If you are using row-based or mixed-format replication, all of the auto-increment lock modes are safe, since row-based replication is not sensitive to the order of execution of the SQL statements (and the mixed format uses row-based replication for any statements that are unsafe for statement-based replication).

- “Lost” auto-increment values and sequence gaps

  In all lock modes (0, 1, and 2), if a transaction that generated auto-increment values rolls back, those auto-increment values are “lost”. Once a value is generated for an auto-increment column, it cannot be rolled back, whether or not the “`INSERT`-like” statement is completed, and whether or not the containing transaction is rolled back. Such lost values are not reused. Thus, there may be gaps in the values stored in an `AUTO_INCREMENT` column of a table.

- Specifying NULL or 0 for the `AUTO_INCREMENT` column

  In all lock modes (0, 1, and 2), if a user specifies NULL or 0 for the `AUTO_INCREMENT` column in an `INSERT`, `InnoDB` treats the row as if the value was not specified and generates a new value for it.

- Assigning a negative value to the `AUTO_INCREMENT` column

  In all lock modes (0, 1, and 2), the behavior of the auto-increment mechanism is not defined if you assign a negative value to the `AUTO_INCREMENT` column.

- If the `AUTO_INCREMENT` value becomes larger than the maximum integer for the specified integer type

  In all lock modes (0, 1, and 2), the behavior of the auto-increment mechanism is not defined if the value becomes larger than the maximum integer that can be stored in the specified integer type.

- Gaps in auto-increment values for “bulk inserts”

  With `innodb_autoinc_lock_mode` set to 0 (“traditional”) or 1 (“consecutive”), the auto-increment values generated by any given statement are consecutive, without gaps, because the table-level `AUTO-INC` lock is held until the end of the statement, and only one such statement can execute at a time.

  With `innodb_autoinc_lock_mode` set to 2 (“interleaved”), there may be gaps in the auto-increment values generated by “bulk inserts,” but only if there are concurrently executing “`INSERT`-like” statements.

  For lock modes 1 or 2, gaps may occur between successive statements because for bulk inserts the exact number of auto-increment values required by each statement may not be known and overestimation is possible.

- Auto-increment values assigned by “mixed-mode inserts”

  Consider a “mixed-mode insert,” where a “simple insert” specifies the auto-increment value for some (but not all) resulting rows. Such a statement behaves differently in lock modes 0, 1, and 2. For example, assume `c1` is an `AUTO_INCREMENT` column of table `t1`, and that the most recent automatically generated sequence number is 100.

```
mysql> CREATE TABLE t1 (
    -> c1 INT UNSIGNED NOT NULL AUTO_INCREMENT PRIMARY KEY,
    -> c2 CHAR(1)
    -> ) ENGINE = INNODB;
```

  Now, consider the following “mixed-mode insert” statement:

```
mysql> INSERT INTO t1 (c1,c2) VALUES (1,'a'), (NULL,'b'), (5,'c'), (NULL,'d');
```

  With `innodb_autoinc_lock_mode` set to 0 (“traditional”), the four new rows are:

```
mysql> SELECT c1, c2 FROM t1 ORDER BY c2;
+-----+------+
| c1  | c2   |
+-----+------+
|   1 | a    |
```

```
| 101 | b    |
|   5 | c    |
| 102 | d    |
+-----+------+
```

The next available auto-increment value is 103 because the auto-increment values are allocated one at a time, not all at once at the beginning of statement execution. This result is true whether or not there are concurrently executing “`INSERT`-like” statements (of any type).

With `innodb_autoinc_lock_mode` set to 1 (“consecutive”), the four new rows are also:

```
mysql> SELECT c1, c2 FROM t1 ORDER BY c2;
+-----+------+
| c1  | c2   |
+-----+------+
|   1 | a    |
| 101 | b    |
|   5 | c    |
| 102 | d    |
+-----+------+
```

However, in this case, the next available auto-increment value is 105, not 103 because four auto-increment values are allocated at the time the statement is processed, but only two are used. This result is true whether or not there are concurrently executing “`INSERT`-like” statements (of any type).

With `innodb_autoinc_lock_mode` set to mode 2 (“interleaved”), the four new rows are:

```
mysql> SELECT c1, c2 FROM t1 ORDER BY c2;
+-----+------+
| c1  | c2   |
+-----+------+
|   1 | a    |
|   x | b    |
|   5 | c    |
|   y | d    |
+-----+------+
```

The values of *x* and *y* are unique and larger than any previously generated rows. However, the specific values of *x* and *y* depend on the number of auto-increment values generated by concurrently executing statements.

Finally, consider the following statement, issued when the most-recently generated sequence number was the value 4:

```
mysql> INSERT INTO t1 (c1,c2) VALUES (1,'a'), (NULL,'b'), (5,'c'), (NULL,'d');
```

With any `innodb_autoinc_lock_mode` setting, this statement generates a duplicate-key error 23000 (`Can't write; duplicate key in table`) because 5 is allocated for the row `(NULL, 'b')` and insertion of the row `(5, 'c')` fails.

- Modifying `AUTO_INCREMENT` column values in the middle of a sequence of `INSERT` statements

  In MySQL 5.7 and earlier, modifying an `AUTO_INCREMENT` column value in the middle of a sequence of `INSERT` statements could lead to “Duplicate entry” errors. For example, if you performed an `UPDATE` operation that changed an `AUTO_INCREMENT` column value to a value larger than the current maximum auto-increment value, subsequent `INSERT` operations that did not specify an unused auto-increment value could encounter “Duplicate entry” errors. In MySQL 8.0 and later, if you modify an `AUTO_INCREMENT` column value to a value larger than the current maximum auto-increment value, the new value is persisted, and subsequent `INSERT` operations allocate auto-increment values starting from the new, larger value. This behavior is demonstrated in the following example.

```
mysql> CREATE TABLE t1 (
    -> c1 INT NOT NULL AUTO_INCREMENT,
    -> PRIMARY KEY (c1)
    ->  ) ENGINE = InnoDB;

mysql> INSERT INTO t1 VALUES(0), (0), (3);

mysql> SELECT c1 FROM t1;
+----+
| c1 |
+----+
|  1 |
|  2 |
|  3 |
+----+

mysql> UPDATE t1 SET c1 = 4 WHERE c1 = 1;

mysql> SELECT c1 FROM t1;
+----+
| c1 |
+----+
|  2 |
|  3 |
|  4 |
+----+

mysql> INSERT INTO t1 VALUES(0);

mysql> SELECT c1 FROM t1;
+----+
| c1 |
+----+
|  2 |
|  3 |
|  4 |
|  5 |
+----+
```

InnoDB AUTO_INCREMENT Counter Initialization

This section describes how `InnoDB` initializes `AUTO_INCREMENT` counters.

If you specify an `AUTO_INCREMENT` column for an `InnoDB` table, the in-memory table object contains a special counter called the auto-increment counter that is used when assigning new values for the column.

In MySQL 5.7 and earlier, the auto-increment counter is stored only in main memory, not on disk. To initialize an auto-increment counter after a server restart, `InnoDB` would execute the equivalent of the following statement on the first insert into a table containing an `AUTO_INCREMENT` column.

```
SELECT MAX(ai_col) FROM table_name FOR UPDATE;
```

In MySQL 8.0, this behavior is changed. The current maximum auto-increment counter value is written to the redo log each time it changes and is saved to an engine-private system table on each checkpoint. These changes make the current maximum auto-increment counter value persistent across server restarts.

On a server restart following a normal shutdown, `InnoDB` initializes the in-memory auto-increment counter using the current maximum auto-increment value stored in the data dictionary system table.

On a server restart during crash recovery, `InnoDB` initializes the in-memory auto-increment counter using the current maximum auto-increment value stored in the data dictionary system table and scans the redo log for auto-increment counter values written since the last checkpoint. If a redo-logged value is greater than the in-memory counter value, the redo-logged value is applied. However, in the case

of a server crash, reuse of a previously allocated auto-increment value cannot be guaranteed. Each time the current maximum auto-increment value is changed due to an `INSERT` or `UPDATE` operation, the new value is written to the redo log, but if the crash occurs before the redo log is flushed to disk, the previously allocated value could be reused when the auto-increment counter is initialized after the server is restarted.

The only circumstance in which `InnoDB` uses the equivalent of a `SELECT MAX(ai_col) FROM table_name FOR UPDATE` statement in MySQL 8.0 and later to initialize an auto-increment counter is when importing a tablespace without a `.cfg` metadata file. Otherwise, the current maximum auto-increment counter value is read from the `.cfg` metadata file.

In MySQL 5.7 and earlier, a server restart cancels the effect of the `AUTO_INCREMENT = N` table option, which may be used in a `CREATE TABLE` or `ALTER TABLE` statement to set an initial counter value or alter the existing counter value, respectively. In MySQL 8.0, a server restart does not cancel the effect of the `AUTO_INCREMENT = N` table option. If you initialize the auto-increment counter to a specific value, or if you alter the auto-increment counter value to a larger value, the new value is persisted across server restarts.

Note

`ALTER TABLE ... AUTO_INCREMENT = N` can only change the auto-increment counter value to a value larger than the current maximum.

In MySQL 5.7 and earlier, a server restart immediately following a `ROLLBACK` operation could result in the reuse of auto-increment values that were previously allocated to the rolled-back transaction, effectively rolling back the current maximum auto-increment value. In MySQL 8.0, the current maximum auto-increment value is persisted, preventing the reuse of previously allocated values.

If a `SHOW TABLE STATUS` statement examines a table before the auto-increment counter is initialized, `InnoDB` opens the table and initializes the counter value using the current maximum auto-increment value that is stored in the data dictionary system table. The value is stored in memory for use by later inserts or updates. Initialization of the counter value uses a normal exclusive-locking read on the table which lasts to the end of the transaction. `InnoDB` follows the same procedure when initializing the auto-increment counter for a newly created table that has a user-specified auto-increment value that is greater than 0.

After the auto-increment counter is initialized, if you do not explicitly specify an auto-increment value when inserting a row, `InnoDB` implicitly increments the counter and assigns the new value to the column. If you insert a row that explicitly specifies an auto-increment column value, and the value is greater than the current maximum counter value, the counter is set to the specified value.

`InnoDB` uses the in-memory auto-increment counter as long as the server runs. When the server is stopped and restarted, `InnoDB` reinitializes the auto-increment counter, as described earlier.

The `auto_increment_offset` configuration option determines the starting point for the `AUTO_INCREMENT` column value. The default setting is 1.

The `auto_increment_increment` configuration option controls the interval between successive column values. The default setting is 1.

15.8.1.6 InnoDB and FOREIGN KEY Constraints

How the `InnoDB` storage engine handles foreign key constraints is described under the following topics in this section:

- Foreign Key Definitions
- Referential Actions
- Foreign Key Restrictions for Generated Columns and Virtual Indexes

- Foreign Key Usage and Error Information

For foreign key usage information and examples, see Section 13.1.15.6, “Using FOREIGN KEY Constraints”.

Foreign Key Definitions

Foreign key definitions for `InnoDB` tables are subject to the following conditions:

- `InnoDB` permits a foreign key to reference any index column or group of columns. However, in the referenced table, there must be an index where the referenced columns are listed as the *first* columns in the same order.
- `InnoDB` does not currently support foreign keys for tables with user-defined partitioning. This means that no user-partitioned `InnoDB` table may contain foreign key references or columns referenced by foreign keys.
- `InnoDB` allows a foreign key constraint to reference a non-unique key. *This is an `InnoDB` extension to standard SQL.*

Referential Actions

Referential actions for foreign keys of `InnoDB` tables are subject to the following conditions:

- While `SET DEFAULT` is allowed by the MySQL Server, it is rejected as invalid by `InnoDB`. `CREATE TABLE` and `ALTER TABLE` statements using this clause are not allowed for InnoDB tables.
- If there are several rows in the parent table that have the same referenced key value, `InnoDB` acts in foreign key checks as if the other parent rows with the same key value do not exist. For example, if you have defined a `RESTRICT` type constraint, and there is a child row with several parent rows, `InnoDB` does not permit the deletion of any of those parent rows.
- `InnoDB` performs cascading operations through a depth-first algorithm, based on records in the indexes corresponding to the foreign key constraints.
- If `ON UPDATE CASCADE` or `ON UPDATE SET NULL` recurses to update the *same table* it has previously updated during the cascade, it acts like `RESTRICT`. This means that you cannot use self-referential `ON UPDATE CASCADE` or `ON UPDATE SET NULL` operations. This is to prevent infinite loops resulting from cascaded updates. A self-referential `ON DELETE SET NULL`, on the other hand, is possible, as is a self-referential `ON DELETE CASCADE`. Cascading operations may not be nested more than 15 levels deep.
- Like MySQL in general, in an SQL statement that inserts, deletes, or updates many rows, `InnoDB` checks `UNIQUE` and `FOREIGN KEY` constraints row-by-row. When performing foreign key checks, `InnoDB` sets shared row-level locks on child or parent records it has to look at. `InnoDB` checks foreign key constraints immediately; the check is not deferred to transaction commit. According to the SQL standard, the default behavior should be deferred checking. That is, constraints are only checked after the *entire SQL statement* has been processed. Until `InnoDB` implements deferred constraint checking, some things are impossible, such as deleting a record that refers to itself using a foreign key.

Foreign Key Restrictions for Generated Columns and Virtual Indexes

- A foreign key constraint on a stored generated column cannot use `ON UPDATE CASCADE`, `ON DELETE SET NULL`, `ON UPDATE SET NULL`, `ON DELETE SET DEFAULT`, or `ON UPDATE SET DEFAULT`.
- A foreign key constraint cannot reference a virtual generated column.
- Prior to MySQL 8.0, a foreign key constraint cannot reference a secondary index defined on a virtual generated column.

Foreign Key Usage and Error Information

You can obtain general information about foreign keys and their usage from querying the `INFORMATION_SCHEMA.KEY_COLUMN_USAGE` table, and more information more specific to `InnoDB` tables can be found in the `INNODB_SYS_FOREIGN` and `INNODB_SYS_FOREIGN_COLS` tables, also in the `INFORMATION_SCHEMA` database.

In addition to `SHOW ERRORS`, in the event of a foreign key error involving `InnoDB` tables (usually Error 150 in the MySQL Server), you can obtain a detailed explanation of the most recent `InnoDB` foreign key error by checking the output of `SHOW ENGINE INNODB STATUS`.

15.8.1.7 Limits on InnoDB Tables

Limits on `InnoDB` tables are described under the following topics in this section:

- Maximums and Minimums
- Restrictions on InnoDB Tables
- Locking and Transactions

Warning

Before using NFS with `InnoDB`, review potential issues outlined in Using NFS with MySQL.

Maximums and Minimums

- A table can contain a maximum of 1017 columns. Virtual generated columns are included in this limit.
- A table can contain a maximum of 64 secondary indexes.
- The index key prefix length limit is 3072 bytes for `InnoDB` tables that use `DYNAMIC` or `COMPRESSED` row format.

  The index key prefix length limit is 767 bytes for `InnoDB` tables that use `REDUNDANT` or `COMPACT` row format. For example, you might hit this limit with a column prefix index of more than 191 characters on a `TEXT` or `VARCHAR` column, assuming a `utf8mb4` character set and the maximum of 4 bytes for each character.

  Attempting to use an index key prefix length that exceeds the limit returns an error.

  The limits that apply to index key prefixes also apply to full-column index keys.
- If you reduce the `InnoDB` page size to 8KB or 4KB by specifying the `innodb_page_size` option when creating the MySQL instance, the maximum length of the index key is lowered proportionally, based on the limit of 3072 bytes for a 16KB page size. That is, the maximum index key length is 1536 bytes when the page size is 8KB, and 768 bytes when the page size is 4KB.
- A maximum of 16 columns is permitted for multicolumn indexes. Exceeding the limit returns an error.

  ```
  ERROR 1070 (42000): Too many key parts specified; max 16 parts allowed
  ```
- The maximum row length, except for variable-length columns (`VARBINARY`, `VARCHAR`, `BLOB` and `TEXT`), is slightly less than half of a page for 4KB, 8KB, 16KB, and 32KB page sizes. For example, the maximum row length for the default `innodb_page_size` of 16KB is about 8000 bytes. For an `InnoDB` page size of 64KB, the maximum row length is about 16000 bytes. `LONGBLOB` and `LONGTEXT` columns must be less than 4GB, and the total row length, including `BLOB` and `TEXT` columns, must be less than 4GB.

  If a row is less than half a page long, all of it is stored locally within the page. If it exceeds half a page, variable-length columns are chosen for external off-page storage until the row fits within half a page, as described in Section 15.11.2, “File Space Management”.

- Although InnoDB supports row sizes larger than 65,535 bytes internally, MySQL itself imposes a row-size limit of 65,535 for the combined size of all columns:

```
mysql> CREATE TABLE t (a VARCHAR(8000), b VARCHAR(10000),
    -> c VARCHAR(10000), d VARCHAR(10000), e VARCHAR(10000),
    -> f VARCHAR(10000), g VARCHAR(10000)) ENGINE=InnoDB;
ERROR 1118 (42000): Row size too large. The maximum row size for the
used table type, not counting BLOBs, is 65535. You have to change some
columns to TEXT or BLOBs
```

  See Section C.10.4, “Limits on Table Column Count and Row Size”.

- On some older operating systems, files must be less than 2GB. This is not a limitation of InnoDB itself, but if you require a large tablespace, configure it using several smaller data files rather than one large data file.

- The combined size of the InnoDB log files can be up to 512GB.

- The minimum tablespace size is slightly larger than 10MB. The maximum tablespace size depends on the InnoDB page size.

  Table 15.6 InnoDB Maximum Tablespace Size

InnoDB Page Size	Maximum Tablespace Size
4KB	16TB
8KB	32TB
16KB	64TB
32KB	128TB
64KB	256TB

  The maximum tablespace size is also the maximum size for a table.

- The default page size in InnoDB is 16KB. You can increase or decrease the page size by configuring the innodb_page_size option when creating the MySQL instance.

  32KB and 64KB page sizes are supported, but ROW_FORMAT=COMPRESSED is unsupported for page sizes greater than 16KB. For both 32KB and 64KB page sizes, the maximum record size is 16KB. For innodb_page_size=32k, extent size is 2MB. For innodb_page_size=64k, extent size is 4MB.

  A MySQL instance using a particular InnoDB page size cannot use data files or log files from an instance that uses a different page size.

Restrictions on InnoDB Tables

- ANALYZE TABLE determines index cardinality (as displayed in the Cardinality column of SHOW INDEX output) by performing random dives on each of the index trees and updating index cardinality estimates accordingly. Because these are only estimates, repeated runs of ANALYZE TABLE could produce different numbers. This makes ANALYZE TABLE fast on InnoDB tables but not 100% accurate because it does not take all rows into account.

  You can make the statistics collected by ANALYZE TABLE more precise and more stable by turning on the innodb_stats_persistent configuration option, as explained in Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”. When that setting is enabled, it is important to run ANALYZE TABLE after major changes to indexed column data, because the statistics are not recalculated periodically (such as after a server restart).

  If the persistent statistics setting is enabled, you can change the number of random dives by modifying the innodb_stats_persistent_sample_pages system variable. If the persistent

statistics setting is disabled, modify the `innodb_stats_transient_sample_pages` system variable instead.

MySQL uses index cardinality estimates in join optimization. If a join is not optimized in the right way, try using `ANALYZE TABLE`. In the few cases that `ANALYZE TABLE` does not produce values good enough for your particular tables, you can use `FORCE INDEX` with your queries to force the use of a particular index, or set the `max_seeks_for_key` system variable to ensure that MySQL prefers index lookups over table scans. See Section B.5.5, “Optimizer-Related Issues”.

- If statements or transactions are running on a table, and `ANALYZE TABLE` is run on the same table followed by a second `ANALYZE TABLE` operation, the second `ANALYZE TABLE` operation is blocked until the statements or transactions are completed. This behavior occurs because `ANALYZE TABLE` marks the currently loaded table definition as obsolete when `ANALYZE TABLE` is finished running. New statements or transactions (including a second `ANALYZE TABLE` statement) must load the new table definition into the table cache, which cannot occur until currently running statements or transactions are completed and the old table definition is purged. Loading multiple concurrent table definitions is not supported.

- `SHOW TABLE STATUS` does not give accurate statistics on `InnoDB` tables except for the physical size reserved by the table. The row count is only a rough estimate used in SQL optimization.

- `InnoDB` does not keep an internal count of rows in a table because concurrent transactions might “see” different numbers of rows at the same time. Consequently, `SELECT COUNT(*)` statements only count rows visible to the current transaction.

  `InnoDB` processes `SELECT COUNT(*)` statements by scanning the clustered index.

  Processing `SELECT COUNT(*)` statements takes some time if index records are not entirely in the buffer pool. For a faster count, you can create a counter table and let your application update it according to the inserts and deletes it does. However, this method may not scale well in situations where thousands of concurrent transactions are initiating updates to the same counter table. If an approximate row count is sufficient, `SHOW TABLE STATUS` can be used.

  `InnoDB` handles `SELECT COUNT(*)` and `SELECT COUNT(1)` operations in the same way. There is no performance difference.

- On Windows, `InnoDB` always stores database and table names internally in lowercase. To move databases in a binary format from Unix to Windows or from Windows to Unix, create all databases and tables using lowercase names.

- An `AUTO_INCREMENT` column *`ai_col`* must be defined as part of an index such that it is possible to perform the equivalent of an indexed `SELECT MAX(`*`ai_col`*`)` lookup on the table to obtain the maximum column value. Typically, this is achieved by making the column the first column of some table index.

- `InnoDB` sets an exclusive lock on the end of the index associated with the `AUTO_INCREMENT` column while initializing a previously specified `AUTO_INCREMENT` column on a table.

  With `innodb_autoinc_lock_mode=0`, `InnoDB` uses a special `AUTO-INC` table lock mode where the lock is obtained and held to the end of the current SQL statement while accessing the auto-increment counter. Other clients cannot insert into the table while the `AUTO-INC` table lock is held. The same behavior occurs for “bulk inserts” with `innodb_autoinc_lock_mode=1`. Table-level `AUTO-INC` locks are not used with `innodb_autoinc_lock_mode=2`. For more information, See Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”.

- When an `AUTO_INCREMENT` integer column runs out of values, a subsequent `INSERT` operation returns a duplicate-key error. This is general MySQL behavior, similar to how `MyISAM` works.

- `DELETE FROM` *`tbl_name`* does not regenerate the table but instead deletes all rows, one by one.

- Cascaded foreign key actions do not activate triggers.

- You cannot create a table with a column name that matches the name of an internal `InnoDB` column (including `DB_ROW_ID`, `DB_TRX_ID`, `DB_ROLL_PTR`, and `DB_MIX_ID`). This restriction applies to use of the names in any letter case.

```
mysql> CREATE TABLE t1 (c1 INT, db_row_id INT) ENGINE=INNODB;
ERROR 1166 (42000): Incorrect column name 'db_row_id'
```

Locking and Transactions

- `LOCK TABLES` acquires two locks on each table if `innodb_table_locks=1` (the default). In addition to a table lock on the MySQL layer, it also acquires an `InnoDB` table lock. Versions of MySQL before 4.1.2 did not acquire `InnoDB` table locks; the old behavior can be selected by setting `innodb_table_locks=0`. If no `InnoDB` table lock is acquired, `LOCK TABLES` completes even if some records of the tables are being locked by other transactions.

  In MySQL 8.0, `innodb_table_locks=0` has no effect for tables locked explicitly with `LOCK TABLES ... WRITE`. It does have an effect for tables locked for read or write by `LOCK TABLES ... WRITE` implicitly (for example, through triggers) or by `LOCK TABLES ... READ`.

- All `InnoDB` locks held by a transaction are released when the transaction is committed or aborted. Thus, it does not make much sense to invoke `LOCK TABLES` on `InnoDB` tables in `autocommit=1` mode because the acquired `InnoDB` table locks would be released immediately.

- You cannot lock additional tables in the middle of a transaction because `LOCK TABLES` performs an implicit `COMMIT` and `UNLOCK TABLES`.

- The limit on data-modifying transactions is 96 * 1023 concurrent transactions that generate undo records. 32 of 128 rollback segments are assigned to non-redo logs for transactions that modify temporary tables and related objects. This means that the maximum number of concurrent data-modifying transactions is 96K. The 96K limit assumes that transactions do not modify temporary tables. If all data-modifying transactions also modify temporary tables, the limit is 32K concurrent transactions.

15.8.2 InnoDB Indexes

This section covers topics related to `InnoDB` indexes.

15.8.2.1 Clustered and Secondary Indexes

Every `InnoDB` table has a special index called the clustered index where the data for the rows is stored. Typically, the clustered index is synonymous with the primary key. To get the best performance from queries, inserts, and other database operations, you must understand how `InnoDB` uses the clustered index to optimize the most common lookup and DML operations for each table.

- When you define a `PRIMARY KEY` on your table, `InnoDB` uses it as the clustered index. Define a primary key for each table that you create. If there is no logical unique and non-null column or set of columns, add a new auto-increment column, whose values are filled in automatically.

- If you do not define a `PRIMARY KEY` for your table, MySQL locates the first `UNIQUE` index where all the key columns are `NOT NULL` and `InnoDB` uses it as the clustered index.

- If the table has no `PRIMARY KEY` or suitable `UNIQUE` index, `InnoDB` internally generates a hidden clustered index on a synthetic column containing row ID values. The rows are ordered by the ID that `InnoDB` assigns to the rows in such a table. The row ID is a 6-byte field that increases monotonically as new rows are inserted. Thus, the rows ordered by the row ID are physically in insertion order.

How the Clustered Index Speeds Up Queries

Accessing a row through the clustered index is fast because the index search leads directly to the page with all the row data. If a table is large, the clustered index architecture often saves a disk I/O operation

when compared to storage organizations that store row data using a different page from the index record. (For example, `MyISAM` uses one file for data rows and another for index records.)

How Secondary Indexes Relate to the Clustered Index

All indexes other than the clustered index are known as secondary indexes. In `InnoDB`, each record in a secondary index contains the primary key columns for the row, as well as the columns specified for the secondary index. `InnoDB` uses this primary key value to search for the row in the clustered index.

If the primary key is long, the secondary indexes use more space, so it is advantageous to have a short primary key.

For guidelines to take advantage of `InnoDB` clustered and secondary indexes, see Section 8.3, “Optimization and Indexes”.

15.8.2.2 The Physical Structure of an InnoDB Index

With the exception of spatial indexes, `InnoDB` indexes are B-tree data structures. Spatial indexes use R-trees, which are specialized data structures for indexing multi-dimensional data. Index records are stored in the leaf pages of their B-tree or R-tree data structure. The default size of an index page is 16KB.

When new records are inserted into an `InnoDB` clustered index, `InnoDB` tries to leave 1/16 of the page free for future insertions and updates of the index records. If index records are inserted in a sequential order (ascending or descending), the resulting index pages are about 15/16 full. If records are inserted in a random order, the pages are from 1/2 to 15/16 full.

`InnoDB` performs a bulk load when creating or rebuilding B-tree indexes. This method of index creation is known as a sorted index build. The `innodb_fill_factor` configuration option defines the percentage of space on each B-tree page that is filled during a sorted index build, with the remaining space reserved for future index growth. Sorted index builds are not supported for spatial indexes. For more information, see Section 15.8.2.3, “Sorted Index Builds”. An `innodb_fill_factor` setting of 100 leaves 1/16 of the space in clustered index pages free for future index growth.

If the fill factor of an `InnoDB` index page drops below the `MERGE_THRESHOLD`, which is 50% by default if not specified, `InnoDB` tries to contract the index tree to free the page. The `MERGE_THRESHOLD` setting applies to both B-tree and R-tree indexes. For more information, see Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

You can define the page size for all `InnoDB` tablespaces in a MySQL instance by setting the `innodb_page_size` configuration option prior to initializing the MySQL instance. Once the page size for an instance is defined, you cannot change it without reinitializing the instance. Supported sizes are 64KB, 32KB, 16KB (default), 8KB, and 4KB, corresponding to the option values `64k`, `32k`, `16k`, `8k`, and `4k`.

A MySQL instance using a particular `InnoDB` page size cannot use data files or log files from an instance that uses a different page size.

15.8.2.3 Sorted Index Builds

`InnoDB` performs a bulk load instead of inserting one index record at a time when creating or rebuilding indexes. This method of index creation is also known as a sorted index build. Sorted index builds are not supported for spatial indexes.

There are three phases to an index build. In the first phase, the clustered index is scanned, and index entries are generated and added to the sort buffer. When the sort buffer becomes full, entries are sorted and written out to a temporary intermediate file. This process is also known as a “run”. In the second phase, with one or more runs written to the temporary intermediate file, a merge sort is performed on all entries in the file. In the third and final phase, the sorted entries are inserted into the B-tree.

Prior to the introduction of sorted index builds, index entries were inserted into the B-tree one record at a time using insert APIs. This method involved opening a B-tree cursor to find the insert position and then inserting entries into a B-tree page using an optimistic insert. If an insert failed due to a page being full, a pessimistic insert would be performed, which involves opening a B-tree cursor and splitting and merging B-tree nodes as necessary to find space for the entry. The drawbacks of this "top-down" method of building an index are the cost of searching for an insert position and the constant splitting and merging of B-tree nodes.

Sorted index builds use a "bottom-up" approach to building an index. With this approach, a reference to the right-most leaf page is held at all levels of the B-tree. The right-most leaf page at the necessary B-tree depth is allocated and entries are inserted according to their sorted order. Once a leaf page is full, a node pointer is appended to the parent page and a sibling leaf page is allocated for the next insert. This process continues until all entries are inserted, which may result in inserts up to the root level. When a sibling page is allocated, the reference to the previously pinned leaf page is released, and the newly allocated leaf page becomes the right-most leaf page and new default insert location.

Reserving B-tree Page Space for Future Index Growth

To set aside space for future index growth, you can use the `innodb_fill_factor` configuration option to reserve a percentage of B-tree page space. For example, setting `innodb_fill_factor` to 80 reserves 20 percent of the space in B-tree pages during a sorted index build. This setting applies to both B-tree leaf and non-leaf pages. It does not apply to external pages used for `TEXT` or `BLOB` entries. The amount of space that is reserved may not be exactly as configured, as the `innodb_fill_factor` value is interpreted as a hint rather than a hard limit.

Sorted Index Builds and Full-Text Index Support

Sorted index builds are supported for fulltext indexes. Previously, SQL was used to insert entries into a fulltext index.

Sorted Index Builds and Compressed Tables

For compressed tables, the previous index creation method appended entries to both compressed and uncompressed pages. When the modification log (representing free space on the compressed page) became full, the compressed page would be recompressed. If compression failed due to a lack of space, the page would be split. With sorted index builds, entries are only appended to uncompressed pages. When an uncompressed page becomes full, it is compressed. Adaptive padding is used to ensure that compression succeeds in most cases, but if compression fails, the page is split and compression is attempted again. This process continues until compression is successful. For more information about compression of B-Tree pages, see Section 15.9.1.5, "How Compression Works for InnoDB Tables".

Sorted Index Builds and Redo Logging

Redo logging is disabled during a sorted index build. Instead, there is a checkpoint to ensure that the index build can withstand a crash or failure. The checkpoint forces a write of all dirty pages to disk. During a sorted index build, the page cleaner thread is signaled periodically to flush dirty pages to ensure that the checkpoint operation can be processed quickly. Normally, the page cleaner thread flushes dirty pages when the number of clean pages falls below a set threshold. For sorted index builds, dirty pages are flushed promptly to reduce checkpoint overhead and to parallelize I/O and CPU activity.

Sorted Index Builds and Optimizer Statistics

Sorted index builds may result in optimizer statistics that differ from those generated by the previous method of index creation. The difference in statistics, which is not expected to affect workload performance, is due to the different algorithm used to populate the index.

15.8.2.4 InnoDB FULLTEXT Indexes

FULLTEXT indexes are created on text-based columns (CHAR, VARCHAR, or TEXT columns) to help speed up queries and DML operations on data contained within those columns, omitting any words that are defined as stopwords.

A FULLTEXT index is defined as part of a CREATE TABLE statement or added to an existing table using ALTER TABLE or CREATE INDEX.

Full-text search is performed using MATCH() ... AGAINST syntax. For usage information, see Section 12.9, “Full-Text Search Functions”.

InnoDB FULLTEXT indexes are described under the following topics in this section:

- InnoDB Full-Text Index Design
- InnoDB Full-Text Index Tables
- InnoDB Full-Text Index Cache
- InnoDB Full-Text Index Document ID and FTS_DOC_ID Column
- InnoDB Full-Text Index Deletion Handling
- InnoDB Full-Text Index Transaction Handling
- Monitoring InnoDB Full-Text Indexes

InnoDB Full-Text Index Design

InnoDB FULLTEXT indexes have an inverted index design. Inverted indexes store a list of words, and for each word, a list of documents that the word appears in. To support proximity search, position information for each word is also stored, as a byte offset.

InnoDB Full-Text Index Tables

For each InnoDB FULLTEXT index, a set of index tables is created, as shown in the following example:

```
mysql> CREATE TABLE opening_lines (
       id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
       opening_line TEXT(500),
       author VARCHAR(200),
       title VARCHAR(200),
       FULLTEXT idx (opening_line)
       ) ENGINE=InnoDB;

mysql> SELECT table_id, name, space from INFORMATION_SCHEMA.INNODB_SYS_TABLES
       WHERE name LIKE 'test/%';
+----------+--------------------------------------------------+-------+
| table_id | name                                             | space |
+----------+--------------------------------------------------+-------+
|      333 | test/fts_0000000000000147_00000000000001c9_index_1 |   289 |
|      334 | test/fts_0000000000000147_00000000000001c9_index_2 |   290 |
|      335 | test/fts_0000000000000147_00000000000001c9_index_3 |   291 |
|      336 | test/fts_0000000000000147_00000000000001c9_index_4 |   292 |
|      337 | test/fts_0000000000000147_00000000000001c9_index_5 |   293 |
|      338 | test/fts_0000000000000147_00000000000001c9_index_6 |   294 |
|      330 | test/fts_0000000000000147_being_deleted            |   286 |
|      331 | test/fts_0000000000000147_being_deleted_cache      |   287 |
|      332 | test/fts_0000000000000147_config                   |   288 |
|      328 | test/fts_0000000000000147_deleted                  |   284 |
|      329 | test/fts_0000000000000147_deleted_cache            |   285 |
|      327 | test/opening_lines                               |   283 |
+----------+--------------------------------------------------+-------+
```

The first six tables represent the inverted index and are referred to as auxiliary index tables. When incoming documents are tokenized, the individual words (also referred to as “tokens”) are inserted into the index tables along with position information and the associated Document ID (DOC_ID). The words

are fully sorted and partitioned among the six index tables based on the character set sort weight of the word's first character.

The inverted index is partitioned into six auxiliary index tables to support parallel index creation. By default, two threads tokenize, sort, and insert words and associated data into the index tables. The number of threads is configurable using the `innodb_ft_sort_pll_degree` option. Consider increasing the number of threads when creating `FULLTEXT` indexes on large tables.

Auxiliary index table names are prefixed with `fts_` and postfixed with `index_*`. Each index table is associated with the indexed table by a hex value in the index table name that matches the `table_id` of the indexed table. For example, the `table_id` of the `test/opening_lines` table is `327`, for which the hex value is 0x147. As shown in the preceding example, the “147” hex value appears in the names of index tables that are associated with the `test/opening_lines` table.

A hex value representing the `index_id` of the `FULLTEXT` index also appears in auxiliary index table names. For example, in the auxiliary table name `test/fts_0000000000000147_00000000000001c9_index_1`, the hex value `1c9` has a decimal value of 457. The index defined on the `opening_lines` table (`idx`) can be identified by querying the `INFORMATION_SCHEMA.INNODB_SYS_INDEXES` table for this value (457).

```
mysql> SELECT index_id, name, table_id, space from INFORMATION_SCHEMA.INNODB_SYS_INDEXES
       WHERE index_id=457;
+----------+------+----------+-------+
| index_id | name | table_id | space |
+----------+------+----------+-------+
|      457 | idx  |      327 |   283 |
+----------+------+----------+-------+
```

Index tables are stored in their own tablespace if the primary table is created in a file-per-table tablespace.

The other index tables shown in the preceding example are used for deletion handling and for storing the internal state of the `FULLTEXT` index.

Auxiliary index table names were changed to lowercase in MySQL 8.0.2.

- `fts_*_deleted` and `fts_*_deleted_cache`

  Contain the document IDs (DOC_ID) for documents that are deleted but whose data is not yet removed from the full-text index. The `fts_*_deleted_cache` is the in-memory version of the `fts_*_deleted` table.

- `fts_*_being_deleted` and `fts_*_being_deleted_cache`

  Contain the document IDs (DOC_ID) for documents that are deleted and whose data is currently in the process of being removed from the full-text index. The `fts_*_being_deleted_cache` table is the in-memory version of the `fts_*_being_deleted` table.

- `fts_*_config`

  Stores information about the internal state of the `FULLTEXT` index. Most importantly, it stores the `FTS_SYNCED_DOC_ID`, which identifies documents that have been parsed and flushed to disk. In case of crash recovery, `FTS_SYNCED_DOC_ID` values are used to identify documents that have not been flushed to disk so that the documents can be re-parsed and added back to the `FULLTEXT` index cache. To view the data in this table, query the `INFORMATION_SCHEMA.INNODB_FT_CONFIG` table.

InnoDB Full-Text Index Cache

When a document is inserted, it is tokenized, and the individual words and associated data are inserted into the `FULLTEXT` index. This process, even for small documents, could result in numerous small insertions into the auxiliary index tables, making concurrent access to these tables a point of

contention. To avoid this problem, InnoDB uses a FULLTEXT index cache to temporarily cache index table insertions for recently inserted rows. This in-memory cache structure holds insertions until the cache is full and then batch flushes them to disk (to the auxiliary index tables). You can query the INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE table to view tokenized data for recently inserted rows.

The caching and batch flushing behavior avoids frequent updates to auxiliary index tables, which could result in concurrent access issues during busy insert and update times. The batching technique also avoids multiple insertions for the same word, and minimizes duplicate entries. Instead of flushing each word individually, insertions for the same word are merged and flushed to disk as a single entry, improving insertion efficiency while keeping auxiliary index tables as small as possible.

The innodb_ft_cache_size variable is used to configure the full-text index cache size (on a per-table basis), which affects how often the full-text index cache is flushed. You can also define a global full-text index cache size limit for all tables in a given instance using the innodb_ft_total_cache_size option.

The full-text index cache stores the same information as auxiliary index tables. However, the full-text index cache only caches tokenized data for recently inserted rows. The data that is already flushed to disk (to the full-text auxiliary tables) is not brought back into the full-text index cache when queried. The data in auxiliary index tables is queried directly, and results from the auxiliary index tables are merged with results from the full-text index cache before being returned.

InnoDB Full-Text Index Document ID and FTS_DOC_ID Column

InnoDB uses a unique document identifier referred to as a Document ID (DOC_ID) to map words in the full-text index to document records where the word appears. The mapping requires an FTS_DOC_ID column on the indexed table. If an FTS_DOC_ID column is not defined, InnoDB automatically adds a hidden FTS_DOC_ID column when the full-text index is created. The following example demonstrates this behavior.

The following table definition does not include an FTS_DOC_ID column:

```
mysql> CREATE TABLE opening_lines (
       id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
       opening_line TEXT(500),
       author VARCHAR(200),
       title VARCHAR(200)
       ) ENGINE=InnoDB;
```

When you create a full-text index on the table using CREATE FULLTEXT INDEX syntax, a warning is returned which reports that InnoDB is rebuilding the table to add the FTS_DOC_ID column.

```
mysql> CREATE FULLTEXT INDEX idx ON opening_lines(opening_line);
Query OK, 0 rows affected, 1 warning (0.19 sec)
Records: 0  Duplicates: 0  Warnings: 1

mysql> SHOW WARNINGS;
+---------+------+-----------------------------------------------------+
| Level   | Code | Message                                             |
+---------+------+-----------------------------------------------------+
| Warning |  124 | InnoDB rebuilding table to add column FTS_DOC_ID |
+---------+------+-----------------------------------------------------+
```

The same warning is returned when using ALTER TABLE to add a full-text index to a table that does not have an FTS_DOC_ID column. If you create a full-text index at CREATE TABLE time and do not specify an FTS_DOC_ID column, InnoDB adds a hidden FTS_DOC_ID column, without warning.

Defining an FTS_DOC_ID column at CREATE TABLE time is less expensive than creating a full-text index on a table that is already loaded with data. If an FTS_DOC_ID column is defined on a table prior to loading data, the table and its indexes do not have to be rebuilt to add the new column. If you are not concerned with CREATE FULLTEXT INDEX performance, leave out the FTS_DOC_ID column to have

InnoDB create it for you. InnoDB creates a hidden FTS_DOC_ID column along with a unique index (FTS_DOC_ID_INDEX) on the FTS_DOC_ID column. If you want to create your own FTS_DOC_ID column, the column must be defined as BIGINT UNSIGNED NOT NULL and named FTS_DOC_ID (all upper case), as in the following example:

Note

The FTS_DOC_ID column does not need to be defined as an AUTO_INCREMENT column, but AUTO_INCREMENT could make loading data easier.

```
mysql> CREATE TABLE opening_lines (
       FTS_DOC_ID BIGINT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
       opening_line TEXT(500),
       author VARCHAR(200),
       title VARCHAR(200)
       ) ENGINE=InnoDB;
```

If you choose to define the FTS_DOC_ID column yourself, you are responsible for managing the column to avoid empty or duplicate values. FTS_DOC_ID values cannot be reused, which means FTS_DOC_ID values must be ever increasing.

Optionally, you can create the required unique FTS_DOC_ID_INDEX (all upper case) on the FTS_DOC_ID column.

```
mysql> CREATE UNIQUE INDEX FTS_DOC_ID_INDEX on opening_lines(FTS_DOC_ID);
```

If you do not create the FTS_DOC_ID_INDEX, InnoDB creates it automatically.

Note

FTS_DOC_ID_INDEX cannot be defined as a descending index because the InnoDB SQL parser does not use descending indexes.

The permitted gap between the largest used FTS_DOC_ID value and new FTS_DOC_ID value is 65535.

InnoDB Full-Text Index Deletion Handling

Deleting a record that has a full-text index column could result in numerous small deletions in the auxiliary index tables, making concurrent access to these tables a point of contention. To avoid this problem, the Document ID (DOC_ID) of a deleted document is logged in a special FTS_*_DELETED table whenever a record is deleted from an indexed table, and the indexed record remains in the full-text index. Before returning query results, information in the FTS_*_DELETED table is used to filter out deleted Document IDs. The benefit of this design is that deletions are fast and inexpensive. The drawback is that the size of the index is not immediately reduced after deleting records. To remove full-text index entries for deleted records, run OPTIMIZE TABLE on the indexed table with innodb_optimize_fulltext_only=ON to rebuild the full-text index. For more information, see Optimizing InnoDB Full-Text Indexes.

InnoDB Full-Text Index Transaction Handling

InnoDB FULLTEXT indexes have special transaction handling characteristics due its caching and batch processing behavior. Specifically, updates and insertions on a FULLTEXT index are processed at transaction commit time, which means that a FULLTEXT search can only see committed data. The following example demonstrates this behavior. The FULLTEXT search only returns a result after the inserted lines are committed.

```
mysql> CREATE TABLE opening_lines (
       id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
       opening_line TEXT(500),
```

```
       author VARCHAR(200),
       title VARCHAR(200),
       FULLTEXT idx (opening_line)
       ) ENGINE=InnoDB;

mysql> BEGIN;

mysql> INSERT INTO opening_lines(opening_line,author,title) VALUES
       ('Call me Ishmael.','Herman Melville','Moby-Dick'),
       ('A screaming comes across the sky.','Thomas Pynchon','Gravity\'s Rainbow'),
       ('I am an invisible man.','Ralph Ellison','Invisible Man'),
       ('Where now? Who now? When now?','Samuel Beckett','The Unnamable'),
       ('It was love at first sight.','Joseph Heller','Catch-22'),
       ('All this happened, more or less.','Kurt Vonnegut','Slaughterhouse-Five'),
       ('Mrs. Dalloway said she would buy the flowers herself.','Virginia Woolf','Mrs. Dalloway'),
       ('It was a pleasure to burn.','Ray Bradbury','Fahrenheit 451');

mysql> SELECT COUNT(*) FROM opening_lines WHERE MATCH(opening_line) AGAINST('Ishmael');
+----------+
| COUNT(*) |
+----------+
|        0 |
+----------+

mysql> COMMIT;

mysql> SELECT COUNT(*) FROM opening_lines WHERE MATCH(opening_line) AGAINST('Ishmael');
+----------+
| COUNT(*) |
+----------+
|        1 |
+----------+
```

Monitoring InnoDB Full-Text Indexes

You can monitor and examine the special text-processing aspects of `InnoDB` `FULLTEXT` indexes by querying the following `INFORMATION_SCHEMA` tables:

- `INNODB_FT_CONFIG`
- `INNODB_FT_INDEX_TABLE`
- `INNODB_FT_INDEX_CACHE`
- `INNODB_FT_DEFAULT_STOPWORD`
- `INNODB_FT_DELETED`
- `INNODB_FT_BEING_DELETED`

You can also view basic information for `FULLTEXT` indexes and tables by querying `INNODB_SYS_INDEXES` and `INNODB_SYS_TABLES`.

For more information, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

15.9 InnoDB Table and Page Compression

This section provides information about the `InnoDB` table compression and `InnoDB` page compression features. The page compression feature is also referred to as transparent page compression.

Using the compression features of `InnoDB`, you can create tables where the data is stored in compressed form. Compression can help to improve both raw performance and scalability. The compression means less data is transferred between disk and memory, and takes up less space on disk and in memory. The benefits are amplified for tables with secondary indexes, because index data is compressed also. Compression can be especially important for SSD storage devices, because they tend to have lower capacity than HDD devices.

15.9.1 InnoDB Table Compression

This section describes `InnoDB` table compression, which is supported with `InnoDB` tables that reside in file_per_table tablespaces or general tablespaces. Table compression is enabled using the `ROW_FORMAT=COMPRESSED` attribute with `CREATE TABLE` or `ALTER TABLE`.

15.9.1.1 Overview of Table Compression

Because processors and cache memories have increased in speed more than disk storage devices, many workloads are disk-bound. Data compression enables smaller database size, reduced I/O, and improved throughput, at the small cost of increased CPU utilization. Compression is especially valuable for read-intensive applications, on systems with enough RAM to keep frequently used data in memory.

An `InnoDB` table created with `ROW_FORMAT=COMPRESSED` can use a smaller page size on disk than the configured `innodb_page_size` value. Smaller pages require less I/O to read from and write to disk, which is especially valuable for SSD devices.

The compressed page size is specified through the `CREATE TABLE` or `ALTER TABLE` `KEY_BLOCK_SIZE` parameter. The different page size requires that the table be placed in a file-per-table tablespace or general tablespace rather than in the system tablespace, as the system tablespace cannot store compressed tables. For more information, see Section 15.7.4, “InnoDB File-Per-Table Tablespaces”, and Section 15.7.9, “InnoDB General Tablespaces”.

The level of compression is the same regardless of the `KEY_BLOCK_SIZE` value. As you specify smaller values for `KEY_BLOCK_SIZE`, you get the I/O benefits of increasingly smaller pages. But if you specify a value that is too small, there is additional overhead to reorganize the pages when data values cannot be compressed enough to fit multiple rows in each page. There is a hard limit on how small `KEY_BLOCK_SIZE` can be for a table, based on the lengths of the key columns for each of its indexes. Specify a value that is too small, and the `CREATE TABLE` or `ALTER TABLE` statement fails.

In the buffer pool, the compressed data is held in small pages, with a page size based on the `KEY_BLOCK_SIZE` value. For extracting or updating the column values, MySQL also creates an uncompressed page in the buffer pool with the uncompressed data. Within the buffer pool, any updates to the uncompressed page are also re-written back to the equivalent compressed page. You might need to size your buffer pool to accommodate the additional data of both compressed and uncompressed pages, although the uncompressed pages are evicted from the buffer pool when space is needed, and then uncompressed again on the next access.

15.9.1.2 Creating Compressed Tables

Compressed tables can be created in file-per-table tablespaces or in general tablespaces. Table compression is not available for the InnoDB system tablespace. The system tablespace (space 0, the .ibdata files) can contain user-created tables, but it also contains internal system data, which is never compressed. Thus, compression applies only to tables (and indexes) stored in file-per-table or general tablespaces.

Creating a Compressed Table in File-Per-Table Tablespace

To create a compressed table in a file-per-table tablespace, `innodb_file_per_table` must be enabled (the default). You can set this parameter in the MySQL configuration file (`my.cnf` or `my.ini`) or dynamically, using a `SET` statement.

After the `innodb_file_per_table` option is configured, specify the `ROW_FORMAT=COMPRESSED` clause or `KEY_BLOCK_SIZE` clause, or both, in a `CREATE TABLE` or `ALTER TABLE` statement to create a compressed table in a file-per-table tablespace.

For example, you might use the following statements:

```
SET GLOBAL innodb_file_per_table=1;
CREATE TABLE t1
```

```
(c1 INT PRIMARY KEY)
ROW_FORMAT=COMPRESSED
KEY_BLOCK_SIZE=8;
```

Creating a Compressed Table in a General Tablespace

To create a compressed table in a general tablespace, `FILE_BLOCK_SIZE` must be defined for the general tablespace, which is specified when the tablespace is created. The `FILE_BLOCK_SIZE` value must be a valid compressed page size in relation to the `innodb_page_size` value, and the page size of the compressed table, defined by the `CREATE TABLE` or `ALTER TABLE KEY_BLOCK_SIZE` clause, must be equal to `FILE_BLOCK_SIZE/1024`. For example, if `innodb_page_size=16384` and `FILE_BLOCK_SIZE=8192`, the `KEY_BLOCK_SIZE` of the table must be 8. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

The following example demonstrates creating a general tablespace and adding a compressed table. The example assumes a default `innodb_page_size` of 16K. The `FILE_BLOCK_SIZE` of 8192 requires that the compressed table have a `KEY_BLOCK_SIZE` of 8.

```
mysql> CREATE TABLESPACE `ts2` ADD DATAFILE 'ts2.ibd' FILE_BLOCK_SIZE = 8192 Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t4 (c1 INT PRIMARY KEY) TABLESPACE ts2 ROW_FORMAT=COMPRESSED KEY_BLOCK_SIZE=8;
Query OK, 0 rows affected (0.00 sec)
```

Notes

- If you specify `ROW_FORMAT=COMPRESSED`, you can omit `KEY_BLOCK_SIZE`; the `KEY_BLOCK_SIZE` setting defaults to half the `innodb_page_size` value.
- If you specify a valid `KEY_BLOCK_SIZE` value, you can omit `ROW_FORMAT=COMPRESSED`; compression is enabled automatically.
- To determine the best value for `KEY_BLOCK_SIZE`, typically you create several copies of the same table with different values for this clause, then measure the size of the resulting `.ibd` files and see how well each performs with a realistic workload. For general tablespaces, keep in mind that dropping a table does not reduce the size of the general tablespace `.ibd` file, nor does it return disk space to the operating system. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.
- The `KEY_BLOCK_SIZE` value is treated as a hint; a different size could be used by `InnoDB` if necessary. For file-per-table tablespaces, the `KEY_BLOCK_SIZE` can only be less than or equal to the `innodb_page_size` value. If you specify a value greater than the `innodb_page_size` value, the specified value is ignored, a warning is issued, and `KEY_BLOCK_SIZE` is set to half of the `innodb_page_size` value. If `innodb_strict_mode=ON`, specifying an invalid `KEY_BLOCK_SIZE` value returns an error. For general tablespaces, valid `KEY_BLOCK_SIZE` values depend on the `FILE_BLOCK_SIZE` setting of the tablespace. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.
- `InnoDB` supports 32k and 64k page sizes but these page sizes do not support compression. For more information, refer to the `innodb_page_size` documentation.
- The default uncompressed size of `InnoDB` data pages is 16KB. Depending on the combination of option values, MySQL uses a page size of 1KB, 2KB, 4KB, 8KB, or 16KB for the tablespace data file (`.ibd` file). The actual compression algorithm is not affected by the `KEY_BLOCK_SIZE` value; the value determines how large each compressed chunk is, which in turn affects how many rows can be packed into each compressed page.
- When creating a compressed table in a file-per-table tablespace, setting `KEY_BLOCK_SIZE` equal to the `InnoDB` page size does not typically result in much compression. For example, setting `KEY_BLOCK_SIZE=16` typically would not result in much compression, since the normal `InnoDB` page size is 16KB. This setting may still be useful for tables with many long `BLOB`, `VARCHAR` or `TEXT` columns, because such values often do compress well, and might therefore require fewer overflow

pages as described in Section 15.9.1.5, “How Compression Works for InnoDB Tables”. For general tablespaces, a KEY_BLOCK_SIZE value equal to the InnoDB page size is not permitted. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

- All indexes of a table (including the clustered index) are compressed using the same page size, as specified in the CREATE TABLE or ALTER TABLE statement. Table attributes such as ROW_FORMAT and KEY_BLOCK_SIZE are not part of the CREATE INDEX syntax for InnoDB tables, and are ignored if they are specified (although, if specified, they will appear in the output of the SHOW CREATE TABLE statement).
- For performance-related configuration options, see Section 15.9.1.3, “Tuning Compression for InnoDB Tables”.

Restrictions on Compressed Tables

- Compressed tables cannot be stored in the InnoDB system tablespace.
- General tablespaces can contain multiple tables, but compressed and uncompressed tables cannot coexist within the same general tablespace.
- Compression applies to an entire table and all its associated indexes, not to individual rows, despite the clause name ROW_FORMAT.
- InnoDB does not support compressed temporary tables. When innodb_strict_mode is enabled (the default), CREATE TEMPORARY TABLE returns errors if ROW_FORMAT=COMPRESSED or KEY_BLOCK_SIZE is specified. If innodb_strict_mode is disabled, warnings are issued and the temporary table is created using a non-compressed row format. The same restrictions apply to ALTER TABLE operations on temporary tables.

15.9.1.3 Tuning Compression for InnoDB Tables

Most often, the internal optimizations described in InnoDB Data Storage and Compression ensure that the system runs well with compressed data. However, because the efficiency of compression depends on the nature of your data, you can make decisions that affect the performance of compressed tables:

- Which tables to compress.
- What compressed page size to use.
- Whether to adjust the size of the buffer pool based on run-time performance characteristics, such as the amount of time the system spends compressing and uncompressing data. Whether the workload is more like a data warehouse (primarily queries) or an OLTP system (mix of queries and DML).
- If the system performs DML operations on compressed tables, and the way the data is distributed leads to expensive compression failures at runtime, you might adjust additional advanced configuration options.

Use the guidelines in this section to help make those architectural and configuration choices. When you are ready to conduct long-term testing and put compressed tables into production, see Section 15.9.1.4, “Monitoring InnoDB Table Compression at Runtime” for ways to verify the effectiveness of those choices under real-world conditions.

When to Use Compression

In general, compression works best on tables that include a reasonable number of character string columns and where the data is read far more often than it is written. Because there are no guaranteed ways to predict whether or not compression benefits a particular situation, always test with a specific workload and data set running on a representative configuration. Consider the following factors when deciding which tables to compress.

Data Characteristics and Compression

A key determinant of the efficiency of compression in reducing the size of data files is the nature of the data itself. Recall that compression works by identifying repeated strings of bytes in a block of data. Completely randomized data is the worst case. Typical data often has repeated values, and so compresses effectively. Character strings often compress well, whether defined in `CHAR`, `VARCHAR`, `TEXT` or `BLOB` columns. On the other hand, tables containing mostly binary data (integers or floating point numbers) or data that is previously compressed (for example JPEG or PNG images) may not generally compress well, significantly or at all.

You choose whether to turn on compression for each InnoDB table. A table and all of its indexes use the same (compressed) page size. It might be that the primary key (clustered) index, which contains the data for all columns of a table, compresses more effectively than the secondary indexes. For those cases where there are long rows, the use of compression might result in long column values being stored “off-page”, as discussed in Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats”. Those overflow pages may compress well. Given these considerations, for many applications, some tables compress more effectively than others, and you might find that your workload performs best only with a subset of tables compressed.

To determine whether or not to compress a particular table, conduct experiments. You can get a rough estimate of how efficiently your data can be compressed by using a utility that implements LZ77 compression (such as `gzip` or WinZip) on a copy of the .ibd file for an uncompressed table. You can expect less compression from a MySQL compressed table than from file-based compression tools, because MySQL compresses data in chunks based on the page size, 16KB by default. In addition to user data, the page format includes some internal system data that is not compressed. File-based compression utilities can examine much larger chunks of data, and so might find more repeated strings in a huge file than MySQL can find in an individual page.

Another way to test compression on a specific table is to copy some data from your uncompressed table to a similar, compressed table (having all the same indexes) in a file-per-table tablespace and look at the size of the resulting `.ibd` file. For example:

```
use test;
set global innodb_file_per_table=1;
set global autocommit=0;

-- Create an uncompressed table with a million or two rows.
create table big_table as select * from information_schema.columns;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
insert into big_table select * from big_table;
commit;
alter table big_table add id int unsigned not null primary key auto_increment;

show create table big_table\G

select count(id) from big_table;

-- Check how much space is needed for the uncompressed table.
\! ls -l data/test/big_table.ibd

create table key_block_size_4 like big_table;
alter table key_block_size_4 key_block_size=4 row_format=compressed;

insert into key_block_size_4 select * from big_table;
commit;

-- Check how much space is needed for a compressed table
-- with particular compression settings.
\! ls -l data/test/key_block_size_4.ibd
```

This experiment produced the following numbers, which of course could vary considerably depending on your table structure and data:

```
-rw-rw----  1 cirrus  staff  310378496 Jan  9 13:44 data/test/big_table.ibd
-rw-rw----  1 cirrus  staff  83886080 Jan  9 15:10 data/test/key_block_size_4.ibd
```

To see whether compression is efficient for your particular workload:

- For simple tests, use a MySQL instance with no other compressed tables and run queries against the `INFORMATION_SCHEMA.INNODB_CMP` table.
- For more elaborate tests involving workloads with multiple compressed tables, run queries against the `INFORMATION_SCHEMA.INNODB_CMP_PER_INDEX` table. Because the statistics in the `INNODB_CMP_PER_INDEX` table are expensive to collect, you must enable the configuration option `innodb_cmp_per_index_enabled` before querying that table, and you might restrict such testing to a development server or a non-critical slave server.
- Run some typical SQL statements against the compressed table you are testing.
- Examine the ratio of successful compression operations to overall compression operations by querying the `INFORMATION_SCHEMA.INNODB_CMP` or `INFORMATION_SCHEMA.INNODB_CMP_PER_INDEX` table, and comparing `COMPRESS_OPS` to `COMPRESS_OPS_OK`.
- If a high percentage of compression operations complete successfully, the table might be a good candidate for compression.
- If you get a high proportion of compression failures, you can adjust `innodb_compression_level`, `innodb_compression_failure_threshold_pct`, and `innodb_compression_pad_pct_max` options as described in Section 15.9.1.6, “Compression for OLTP Workloads”, and try further tests.

Database Compression versus Application Compression

Decide whether to compress data in your application or in the table; do not use both types of compression for the same data. When you compress the data in the application and store the results in a compressed table, extra space savings are extremely unlikely, and the double compression just wastes CPU cycles.

Compressing in the Database

When enabled, MySQL table compression is automatic and applies to all columns and index values. The columns can still be tested with operators such as `LIKE`, and sort operations can still use indexes even when the index values are compressed. Because indexes are often a significant fraction of the total size of a database, compression could result in significant savings in storage, I/O or processor time. The compression and decompression operations happen on the database server, which likely is a powerful system that is sized to handle the expected load.

Compressing in the Application

If you compress data such as text in your application, before it is inserted into the database, You might save overhead for data that does not compress well by compressing some columns and not others. This approach uses CPU cycles for compression and uncompression on the client machine rather than the database server, which might be appropriate for a distributed application with many clients, or where the client machine has spare CPU cycles.

Hybrid Approach

Of course, it is possible to combine these approaches. For some applications, it may be appropriate to use some compressed tables and some uncompressed tables. It may be best to externally compress some data (and store it in uncompressed tables) and allow MySQL to compress (some of) the other tables in the application. As always, up-front design and real-life testing are valuable in reaching the right decision.

Workload Characteristics and Compression

In addition to choosing which tables to compress (and the page size), the workload is another key determinant of performance. If the application is dominated by reads, rather than updates, fewer pages need to be reorganized and recompressed after the index page runs out of room for the per-page "modification log" that MySQL maintains for compressed data. If the updates predominantly change non-indexed columns or those containing `BLOB`s or large strings that happen to be stored "off-page", the overhead of compression may be acceptable. If the only changes to a table are `INSERT`s that use a monotonically increasing primary key, and there are few secondary indexes, there is little need to reorganize and recompress index pages. Since MySQL can "delete-mark" and delete rows on compressed pages "in place" by modifying uncompressed data, `DELETE` operations on a table are relatively efficient.

For some environments, the time it takes to load data can be as important as run-time retrieval. Especially in data warehouse environments, many tables may be read-only or read-mostly. In those cases, it might or might not be acceptable to pay the price of compression in terms of increased load time, unless the resulting savings in fewer disk reads or in storage cost is significant.

Fundamentally, compression works best when the CPU time is available for compressing and uncompressing data. Thus, if your workload is I/O bound, rather than CPU-bound, you might find that compression can improve overall performance. When you test your application performance with different compression configurations, test on a platform similar to the planned configuration of the production system.

Configuration Characteristics and Compression

Reading and writing database pages from and to disk is the slowest aspect of system performance. Compression attempts to reduce I/O by using CPU time to compress and uncompress data, and is most effective when I/O is a relatively scarce resource compared to processor cycles.

This is often especially the case when running in a multi-user environment with fast, multi-core CPUs. When a page of a compressed table is in memory, MySQL often uses additional memory, typically 16KB, in the buffer pool for an uncompressed copy of the page. The adaptive LRU algorithm attempts to balance the use of memory between compressed and uncompressed pages to take into account whether the workload is running in an I/O-bound or CPU-bound manner. Still, a configuration with more memory dedicated to the buffer pool tends to run better when using compressed tables than a configuration where memory is highly constrained.

Choosing the Compressed Page Size

The optimal setting of the compressed page size depends on the type and distribution of data that the table and its indexes contain. The compressed page size should always be bigger than the maximum record size, or operations may fail as noted in Compression of B-Tree Pages.

Setting the compressed page size too large wastes some space, but the pages do not have to be compressed as often. If the compressed page size is set too small, inserts or updates may require time-consuming recompression, and the B-tree nodes may have to be split more frequently, leading to bigger data files and less efficient indexing.

Typically, you set the compressed page size to 8K or 4K bytes. Given that the maximum row size for an InnoDB table is around 8K, `KEY_BLOCK_SIZE=8` is usually a safe choice.

15.9.1.4 Monitoring InnoDB Table Compression at Runtime

Overall application performance, CPU and I/O utilization and the size of disk files are good indicators of how effective compression is for your application. This section builds on the performance tuning advice from Section 15.9.1.3, "Tuning Compression for InnoDB Tables", and shows how to find problems that might not turn up during initial testing.

To dig deeper into performance considerations for compressed tables, you can monitor compression performance at runtime using the Information Schema tables described in Example 15.1, "Using the

Compression Information Schema Tables". These tables reflect the internal use of memory and the rates of compression used overall.

The `INNODB_CMP` table reports information about compression activity for each compressed page size (`KEY_BLOCK_SIZE`) in use. The information in these tables is system-wide: it summarizes the compression statistics across all compressed tables in your database. You can use this data to help decide whether or not to compress a table by examining these tables when no other compressed tables are being accessed. It involves relatively low overhead on the server, so you might query it periodically on a production server to check the overall efficiency of the compression feature.

The `INNODB_CMP_PER_INDEX` table reports information about compression activity for individual tables and indexes. This information is more targeted and more useful for evaluating compression efficiency and diagnosing performance issues one table or index at a time. (Because that each `InnoDB` table is represented as a clustered index, MySQL does not make a big distinction between tables and indexes in this context.) The `INNODB_CMP_PER_INDEX` table does involve substantial overhead, so it is more suitable for development servers, where you can compare the effects of different workloads, data, and compression settings in isolation. To guard against imposing this monitoring overhead by accident, you must enable the `innodb_cmp_per_index_enabled` configuration option before you can query the `INNODB_CMP_PER_INDEX` table.

The key statistics to consider are the number of, and amount of time spent performing, compression and uncompression operations. Since MySQL splits B-tree nodes when they are too full to contain the compressed data following a modification, compare the number of "successful" compression operations with the number of such operations overall. Based on the information in the `INNODB_CMP` and `INNODB_CMP_PER_INDEX` tables and overall application performance and hardware resource utilization, you might make changes in your hardware configuration, adjust the size of the buffer pool, choose a different page size, or select a different set of tables to compress.

If the amount of CPU time required for compressing and uncompressing is high, changing to faster or multi-core CPUs can help improve performance with the same data, application workload and set of compressed tables. Increasing the size of the buffer pool might also help performance, so that more uncompressed pages can stay in memory, reducing the need to uncompress pages that exist in memory only in compressed form.

A large number of compression operations overall (compared to the number of `INSERT`, `UPDATE` and `DELETE` operations in your application and the size of the database) could indicate that some of your compressed tables are being updated too heavily for effective compression. If so, choose a larger page size, or be more selective about which tables you compress.

If the number of "successful" compression operations (`COMPRESS_OPS_OK`) is a high percentage of the total number of compression operations (`COMPRESS_OPS`), then the system is likely performing well. If the ratio is low, then MySQL is reorganizing, recompressing, and splitting B-tree nodes more often than is desirable. In this case, avoid compressing some tables, or increase `KEY_BLOCK_SIZE` for some of the compressed tables. You might turn off compression for tables that cause the number of "compression failures" in your application to be more than 1% or 2% of the total. (Such a failure ratio might be acceptable during a temporary operation such as a data load).

15.9.1.5 How Compression Works for InnoDB Tables

This section describes some internal implementation details about compression for InnoDB tables. The information presented here may be helpful in tuning for performance, but is not necessary to know for basic use of compression.

Compression Algorithms

Some operating systems implement compression at the file system level. Files are typically divided into fixed-size blocks that are compressed into variable-size blocks, which easily leads into fragmentation. Every time something inside a block is modified, the whole block is recompressed before it is written to disk. These properties make this compression technique unsuitable for use in an update-intensive database system.

MySQL implements compression with the help of the well-known zlib library, which implements the LZ77 compression algorithm. This compression algorithm is mature, robust, and efficient in both CPU utilization and in reduction of data size. The algorithm is “lossless”, so that the original uncompressed data can always be reconstructed from the compressed form. LZ77 compression works by finding sequences of data that are repeated within the data to be compressed. The patterns of values in your data determine how well it compresses, but typical user data often compresses by 50% or more.

Unlike compression performed by an application, or compression features of some other database management systems, InnoDB compression applies both to user data and to indexes. In many cases, indexes can constitute 40-50% or more of the total database size, so this difference is significant. When compression is working well for a data set, the size of the InnoDB data files (the file-per-table tablespace or general tablespace `.idb` files) is 25% to 50% of the uncompressed size or possibly smaller. Depending on the workload, this smaller database can in turn lead to a reduction in I/O, and an increase in throughput, at a modest cost in terms of increased CPU utilization. You can adjust the balance between compression level and CPU overhead by modifying the `innodb_compression_level` configuration option.

InnoDB Data Storage and Compression

All user data in InnoDB tables is stored in pages comprising a B-tree index (the clustered index). In some other database systems, this type of index is called an “index-organized table”. Each row in the index node contains the values of the (user-specified or system-generated) primary key and all the other columns of the table.

Secondary indexes in InnoDB tables are also B-trees, containing pairs of values: the index key and a pointer to a row in the clustered index. The pointer is in fact the value of the primary key of the table, which is used to access the clustered index if columns other than the index key and primary key are required. Secondary index records must always fit on a single B-tree page.

The compression of B-tree nodes (of both clustered and secondary indexes) is handled differently from compression of overflow pages used to store long `VARCHAR`, `BLOB`, or `TEXT` columns, as explained in the following sections.

Compression of B-Tree Pages

Because they are frequently updated, B-tree pages require special treatment. It is important to minimize the number of times B-tree nodes are split, as well as to minimize the need to uncompress and recompress their content.

One technique MySQL uses is to maintain some system information in the B-tree node in uncompressed form, thus facilitating certain in-place updates. For example, this allows rows to be delete-marked and deleted without any compression operation.

In addition, MySQL attempts to avoid unnecessary uncompression and recompression of index pages when they are changed. Within each B-tree page, the system keeps an uncompressed “modification log” to record changes made to the page. Updates and inserts of small records may be written to this modification log without requiring the entire page to be completely reconstructed.

When the space for the modification log runs out, InnoDB uncompresses the page, applies the changes and recompresses the page. If recompression fails (a situation known as a compression failure), the B-tree nodes are split and the process is repeated until the update or insert succeeds.

To avoid frequent compression failures in write-intensive workloads, such as for OLTP applications, MySQL sometimes reserves some empty space (padding) in the page, so that the modification log fills up sooner and the page is recompressed while there is still enough room to avoid splitting it. The amount of padding space left in each page varies as the system keeps track of the frequency of page splits. On a busy server doing frequent writes to compressed tables, you can adjust the `innodb_compression_failure_threshold_pct`, and `innodb_compression_pad_pct_max` configuration options to fine-tune this mechanism.

Generally, MySQL requires that each B-tree page in an InnoDB table can accommodate at least two records. For compressed tables, this requirement has been relaxed. Leaf pages of B-tree nodes (whether of the primary key or secondary indexes) only need to accommodate one record, but that record must fit, in uncompressed form, in the per-page modification log. If `innodb_strict_mode` is `ON`, MySQL checks the maximum row size during `CREATE TABLE` or `CREATE INDEX`. If the row does not fit, the following error message is issued: `ERROR HY000: Too big row`.

If you create a table when `innodb_strict_mode` is OFF, and a subsequent `INSERT` or `UPDATE` statement attempts to create an index entry that does not fit in the size of the compressed page, the operation fails with `ERROR 42000: Row size too large`. (This error message does not name the index for which the record is too large, or mention the length of the index record or the maximum record size on that particular index page.) To solve this problem, rebuild the table with `ALTER TABLE` and select a larger compressed page size (`KEY_BLOCK_SIZE`), shorten any column prefix indexes, or disable compression entirely with `ROW_FORMAT=DYNAMIC` or `ROW_FORMAT=COMPACT`.

`innodb_strict_mode` is not applicable to general tablespaces, which also support compressed tables. Tablespace management rules for general tablespaces are strictly enforced independently of `innodb_strict_mode`. For more information, see Section 13.1.16, “CREATE TABLESPACE Syntax”.

Compressing BLOB, VARCHAR, and TEXT Columns

In an InnoDB table, `BLOB`, `VARCHAR`, and `TEXT` columns that are not part of the primary key may be stored on separately allocated overflow pages. We refer to these columns as off-page columns. Their values are stored on singly-linked lists of overflow pages.

For tables created in `ROW_FORMAT=DYNAMIC` or `ROW_FORMAT=COMPRESSED`, the values of `BLOB`, `TEXT`, or `VARCHAR` columns may be stored fully off-page, depending on their length and the length of the entire row. For columns that are stored off-page, the clustered index record only contains 20-byte pointers to the overflow pages, one per column. Whether any columns are stored off-page depends on the page size and the total size of the row. When the row is too long to fit entirely within the page of the clustered index, MySQL chooses the longest columns for off-page storage until the row fits on the clustered index page. As noted above, if a row does not fit by itself on a compressed page, an error occurs.

Note

For tables created in `ROW_FORMAT=DYNAMIC` or `ROW_FORMAT=COMPRESSED`, `TEXT` and `BLOB` columns that are less than or equal to 40 bytes are always stored in-line.

Tables that use `ROW_FORMAT=REDUNDANT` and `ROW_FORMAT=COMPACT` store the first 768 bytes of `BLOB`, `VARCHAR`, and `TEXT` columns in the clustered index record along with the primary key. The 768-byte prefix is followed by a 20-byte pointer to the overflow pages that contain the rest of the column value.

When a table is in `COMPRESSED` format, all data written to overflow pages is compressed “as is”; that is, MySQL applies the zlib compression algorithm to the entire data item. Other than the data, compressed overflow pages contain an uncompressed header and trailer comprising a page checksum and a link to the next overflow page, among other things. Therefore, very significant storage savings can be obtained for longer `BLOB`, `TEXT`, or `VARCHAR` columns if the data is highly compressible, as is often the case with text data. Image data, such as `JPEG`, is typically already compressed and so does not benefit much from being stored in a compressed table; the double compression can waste CPU cycles for little or no space savings.

The overflow pages are of the same size as other pages. A row containing ten columns stored off-page occupies ten overflow pages, even if the total length of the columns is only 8K bytes. In an uncompressed table, ten uncompressed overflow pages occupy 160K bytes. In a compressed table with an 8K page size, they occupy only 80K bytes. Thus, it is often more efficient to use compressed table format for tables with long column values.

For file-per-table tablespaces, using a 16K compressed page size can reduce storage and I/O costs for `BLOB`, `VARCHAR`, or `TEXT` columns, because such data often compress well, and might therefore require fewer overflow pages, even though the B-tree nodes themselves take as many pages as in the uncompressed form. General tablespaces do not support a 16K compressed page size (`KEY_BLOCK_SIZE`). For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

Compression and the InnoDB Buffer Pool

In a compressed `InnoDB` table, every compressed page (whether 1K, 2K, 4K or 8K) corresponds to an uncompressed page of 16K bytes (or a smaller size if `innodb_page_size` is set). To access the data in a page, MySQL reads the compressed page from disk if it is not already in the buffer pool, then uncompresses the page to its original form. This section describes how `InnoDB` manages the buffer pool with respect to pages of compressed tables.

To minimize I/O and to reduce the need to uncompress a page, at times the buffer pool contains both the compressed and uncompressed form of a database page. To make room for other required database pages, MySQL can evict from the buffer pool an uncompressed page, while leaving the compressed page in memory. Or, if a page has not been accessed in a while, the compressed form of the page might be written to disk, to free space for other data. Thus, at any given time, the buffer pool might contain both the compressed and uncompressed forms of the page, or only the compressed form of the page, or neither.

MySQL keeps track of which pages to keep in memory and which to evict using a least-recently-used (LRU) list, so that hot (frequently accessed) data tends to stay in memory. When compressed tables are accessed, MySQL uses an adaptive LRU algorithm to achieve an appropriate balance of compressed and uncompressed pages in memory. This adaptive algorithm is sensitive to whether the system is running in an I/O-bound or CPU-bound manner. The goal is to avoid spending too much processing time uncompressing pages when the CPU is busy, and to avoid doing excess I/O when the CPU has spare cycles that can be used for uncompressing compressed pages (that may already be in memory). When the system is I/O-bound, the algorithm prefers to evict the uncompressed copy of a page rather than both copies, to make more room for other disk pages to become memory resident. When the system is CPU-bound, MySQL prefers to evict both the compressed and uncompressed page, so that more memory can be used for “hot” pages and reducing the need to uncompress data in memory only in compressed form.

Compression and the InnoDB Redo Log Files

Before a compressed page is written to a data file, MySQL writes a copy of the page to the redo log (if it has been recompressed since the last time it was written to the database). This is done to ensure that redo logs are usable for crash recovery, even in the unlikely case that the `zlib` library is upgraded and that change introduces a compatibility problem with the compressed data. Therefore, some increase in the size of log files, or a need for more frequent checkpoints, can be expected when using compression. The amount of increase in the log file size or checkpoint frequency depends on the number of times compressed pages are modified in a way that requires reorganization and recompression.

To create a compressed table in a file-per-table tablespace, `innodb_file_per_table` must be enabled. There is no dependence on the `innodb_file_per_table` setting when creating a compressed table in a general tablespace. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

15.9.1.6 Compression for OLTP Workloads

Traditionally, the `InnoDB` compression feature was recommended primarily for read-only or read-mostly workloads, such as in a data warehouse configuration. The rise of SSD storage devices, which are fast but relatively small and expensive, makes compression attractive also for `OLTP` workloads: high-traffic, interactive web sites can reduce their storage requirements and their I/O operations per second (IOPS) by using compressed tables with applications that do frequent `INSERT`, `UPDATE`, and `DELETE` operations.

These configuration options let you adjust the way compression works for a particular MySQL instance, with an emphasis on performance and scalability for write-intensive operations:

- `innodb_compression_level` lets you turn the degree of compression up or down. A higher value lets you fit more data onto a storage device, at the expense of more CPU overhead during compression. A lower value lets you reduce CPU overhead when storage space is not critical, or you expect the data is not especially compressible.
- `innodb_compression_failure_threshold_pct` specifies a cutoff point for compression failures during updates to a compressed table. When this threshold is passed, MySQL begins to leave additional free space within each new compressed page, dynamically adjusting the amount of free space up to the percentage of page size specified by `innodb_compression_pad_pct_max`
- `innodb_compression_pad_pct_max` lets you adjust the maximum amount of space reserved within each page to record changes to compressed rows, without needing to compress the entire page again. The higher the value, the more changes can be recorded without recompressing the page. MySQL uses a variable amount of free space for the pages within each compressed table, only when a designated percentage of compression operations "fail" at runtime, requiring an expensive operation to split the compressed page.
- `innodb_log_compressed_pages` lets you disable writing of images of re-compressed pages to the redo log. Re-compression may occur when changes are made to compressed data. This option is enabled by default to prevent corruption that could occur if a different version of the `zlib` compression algorithm is used during recovery. If you are certain that the `zlib` version will not change, disable `innodb_log_compressed_pages` to reduce redo log generation for workloads that modify compressed data.

Because working with compressed data sometimes involves keeping both compressed and uncompressed versions of a page in memory at the same time, when using compression with an OLTP-style workload, be prepared to increase the value of the `innodb_buffer_pool_size` configuration option.

15.9.1.7 SQL Compression Syntax Warnings and Errors

This section describes syntax warnings and errors that you may encounter when using the table compression feature with file-per-table tablespaces and general tablespaces.

SQL Compression Syntax Warnings and Errors for File-Per-Table Tablespaces

When `innodb_strict_mode` is enabled (the default), specifying `ROW_FORMAT=COMPRESSED` or `KEY_BLOCK_SIZE` in `CREATE TABLE` or `ALTER TABLE` statements produces the following error if `innodb_file_per_table` is disabled.

```
ERROR 1031 (HY000): Table storage engine for 't1' doesn't have this option
```

Note

The table is not created if the current configuration does not permit using compressed tables.

When `innodb_strict_mode` is disabled, specifying `ROW_FORMAT=COMPRESSED` or `KEY_BLOCK_SIZE` in `CREATE TABLE` or `ALTER TABLE` statements produces the following warnings if `innodb_file_per_table` is disabled.

```
mysql> SHOW WARNINGS;
+---------+------+-----------------------------------------------------------------+
| Level   | Code | Message                                                         |
+---------+------+-----------------------------------------------------------------+
| Warning | 1478 | InnoDB: KEY_BLOCK_SIZE requires innodb_file_per_table.          |
| Warning | 1478 | InnoDB: ignoring KEY_BLOCK_SIZE=4.                              |
| Warning | 1478 | InnoDB: ROW_FORMAT=COMPRESSED requires innodb_file_per_table. |
| Warning | 1478 | InnoDB: assuming ROW_FORMAT=DYNAMIC.                            |
+---------+------+-----------------------------------------------------------------+
```

Note

These messages are only warnings, not errors, and the table is created without compression, as if the options were not specified.

The “non-strict” behavior lets you import a `mysqldump` file into a database that does not support compressed tables, even if the source database contained compressed tables. In that case, MySQL creates the table in `ROW_FORMAT=DYNAMIC` instead of preventing the operation.

To import the dump file into a new database, and have the tables re-created as they exist in the original database, ensure the server has the proper setting for the `innodb_file_per_table` configuration parameter.

The attribute `KEY_BLOCK_SIZE` is permitted only when `ROW_FORMAT` is specified as `COMPRESSED` or is omitted. Specifying a `KEY_BLOCK_SIZE` with any other `ROW_FORMAT` generates a warning that you can view with `SHOW WARNINGS`. However, the table is non-compressed; the specified `KEY_BLOCK_SIZE` is ignored).

Level	Code	Message
Warning	1478	`InnoDB: ignoring KEY_BLOCK_SIZE=`*`n`*` unless ROW_FORMAT=COMPRESSED.`

If you are running with `innodb_strict_mode` enabled, the combination of a `KEY_BLOCK_SIZE` with any `ROW_FORMAT` other than `COMPRESSED` generates an error, not a warning, and the table is not created.

Table 15.7, “ROW_FORMAT and KEY_BLOCK_SIZE Options” provides an overview the `ROW_FORMAT` and `KEY_BLOCK_SIZE` options that are used with `CREATE TABLE` or `ALTER TABLE`.

Table 15.7 ROW_FORMAT and KEY_BLOCK_SIZE Options

Option	Usage Notes	Description
`ROW_FORMAT=REDUNDANT`	Storage format used prior to MySQL 5.0.3	Less efficient than `ROW_FORMAT=COMPACT`; for backward compatibility
`ROW_FORMAT=COMPACT`	Default storage format since MySQL 5.0.3	Stores a prefix of 768 bytes of long column values in the clustered index page, with the remaining bytes stored in an overflow page
`ROW_FORMAT=DYNAMIC`		Store values within the clustered index page if they fit; if not, stores only a 20-byte pointer to an overflow page (no prefix)
`ROW_FORMAT=COMPRESSED`		Compresses the table and indexes using zlib
`KEY_BLOCK_SIZE=`*`n`		Specifies compressed page size of 1, 2, 4, 8 or 16 kilobytes; implies `ROW_FORMAT=COMPRESSED`. For general tablespaces, a `KEY_BLOCK_SIZE` value equal to the `InnoDB` page size is not permitted.

Table 15.8, “CREATE/ALTER TABLE Warnings and Errors when InnoDB Strict Mode is OFF” summarizes error conditions that occur with certain combinations of configuration parameters and options on the `CREATE TABLE` or `ALTER TABLE` statements, and how the options appear in the output of `SHOW TABLE STATUS`.

When `innodb_strict_mode` is `OFF`, MySQL creates or alters the table, but ignores certain settings as shown below. You can see the warning messages in the MySQL error log. When `innodb_strict_mode` is `ON`, these specified combinations of options generate errors, and the table is not created or altered. To see the full description of the error condition, issue the `SHOW ERRORS` statement: example:

```
mysql> CREATE TABLE x (id INT PRIMARY KEY, c INT)

-> ENGINE=INNODB KEY_BLOCK_SIZE=33333;

ERROR 1005 (HY000): Can't create table 'test.x' (errno: 1478)

mysql> SHOW ERRORS;
+-------+------+-----------------------------------------+
| Level | Code | Message                                 |
+-------+------+-----------------------------------------+
| Error | 1478 | InnoDB: invalid KEY_BLOCK_SIZE=33333.   |
| Error | 1005 | Can't create table 'test.x' (errno: 1478) |
+-------+------+-----------------------------------------+
```

Table 15.8 CREATE/ALTER TABLE Warnings and Errors when InnoDB Strict Mode is OFF

Syntax	Warning or Error Condition	Resulting `ROW_FORMAT`, as shown in `SHOW TABLE STATUS`
`ROW_FORMAT=REDUNDANT`	None	`REDUNDANT`
`ROW_FORMAT=COMPACT`	None	`COMPACT`
`ROW_FORMAT=COMPRESSED` or `ROW_FORMAT=DYNAMIC` or `KEY_BLOCK_SIZE` is specified	Ignored for file-per-table tablespaces unless `innodb_file_per_table` is enabled. General tablespaces support all row formats. See Section 15.7.9, "InnoDB General Tablespaces".	`the default row format for file-per-table tablespaces; the specified row format for general tablespaces`
Invalid `KEY_BLOCK_SIZE` is specified (not 1, 2, 4, 8 or 16)	`KEY_BLOCK_SIZE` is ignored	the specified row format, or the default row format
`ROW_FORMAT=COMPRESSED` and valid `KEY_BLOCK_SIZE` are specified	None; `KEY_BLOCK_SIZE` specified is used	`COMPRESSED`
`KEY_BLOCK_SIZE` is specified with `REDUNDANT`, `COMPACT` or `DYNAMIC` row format	`KEY_BLOCK_SIZE` is ignored	`REDUNDANT`, `COMPACT` or `DYNAMIC`
`ROW_FORMAT` is not one of `REDUNDANT`, `COMPACT`, `DYNAMIC` or `COMPRESSED`	Ignored if recognized by the MySQL parser. Otherwise, an error is issued.	the default row format or N/A

When `innodb_strict_mode` is `ON`, MySQL rejects invalid `ROW_FORMAT` or `KEY_BLOCK_SIZE` parameters and issues errors. Strict mode is `ON` by default. When `innodb_strict_mode` is `OFF`, MySQL issues warnings instead of errors for ignored invalid parameters.

It is not possible to see the chosen `KEY_BLOCK_SIZE` using `SHOW TABLE STATUS`. The statement `SHOW CREATE TABLE` displays the `KEY_BLOCK_SIZE` (even if it was ignored when creating the table). The real compressed page size of the table cannot be displayed by MySQL.

SQL Compression Syntax Warnings and Errors for General Tablespaces

- If `FILE_BLOCK_SIZE` was not defined for the general tablespace when the tablespace was created, the tablespace cannot contain compressed tables. If you attempt to add a compressed table, an error is returned, as shown in the following example:

```
mysql> CREATE TABLESPACE `ts1` ADD DATAFILE 'ts1.ibd' Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY) TABLESPACE ts1 ROW_FORMAT=COMPRESSED
```

```
KEY_BLOCK_SIZE=8;
ERROR 1478 (HY000): InnoDB: Tablespace `ts1` cannot contain a COMPRESSED table
```

- Attempting to add a table with an invalid KEY_BLOCK_SIZE to a general tablespace returns an error, as shown in the following example:

```
mysql> CREATE TABLESPACE `ts2` ADD DATAFILE 'ts2.ibd' FILE_BLOCK_SIZE = 8192 Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t2 (c1 INT PRIMARY KEY) TABLESPACE ts2 ROW_FORMAT=COMPRESSED
KEY_BLOCK_SIZE=4;
ERROR 1478 (HY000): InnoDB: Tablespace `ts2` uses block size 8192 and cannot
contain a table with physical page size 4096
```

  For general tablespaces, the KEY_BLOCK_SIZE of the table must be equal to the FILE_BLOCK_SIZE of the tablespace divided by 1024. For example, if the FILE_BLOCK_SIZE of the tablespace is 8192, the KEY_BLOCK_SIZE of the table must be 8.

- Attempting to add a table with an uncompressed row format to a general tablespace configured to store compressed tables returns an error, as shown in the following example:

```
mysql> CREATE TABLESPACE `ts3` ADD DATAFILE 'ts3.ibd' FILE_BLOCK_SIZE = 8192 Engine=InnoDB;
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE t3 (c1 INT PRIMARY KEY) TABLESPACE ts3 ROW_FORMAT=COMPACT;
ERROR 1478 (HY000): InnoDB: Tablespace `ts3` uses block size 8192 and cannot
contain a table with physical page size 16384
```

innodb_strict_mode is not applicable to general tablespaces. Tablespace management rules for general tablespaces are strictly enforced independently of innodb_strict_mode. For more information, see Section 13.1.16, “CREATE TABLESPACE Syntax”.

For more information about using compressed tables with general tablespaces, see Section 15.7.9, “InnoDB General Tablespaces”.

15.9.2 InnoDB Page Compression

InnoDB supports page-level compression for tables that reside in file-per-table tablespaces. This feature is referred to as *Transparent Page Compression*. Page compression is enabled by specifying the COMPRESSION attribute with CREATE TABLE or ALTER TABLE. Supported compression algorithms include Zlib and LZ4.

Supported Platforms

Page compression requires sparse file and hole punching support. Page compression is supported on Windows with NTFS, and on the following subset of MySQL-supported Linux platforms where the kernel level provides hole punching support:

- RHEL 7 and derived distributions that use kernel version 3.10.0-123 or higher
- OEL 5.10 (UEK2) kernel version 2.6.39 or higher
- OEL 6.5 (UEK3) kernel version 3.8.13 or higher
- OEL 7.0 kernel version 3.8.13 or higher
- SLE11 kernel version 3.0-x
- SLE12 kernel version 3.12-x
- OES11 kernel version 3.0-x
- Ubuntu 14.0.4 LTS kernel version 3.13 or higher
- Ubuntu 12.0.4 LTS kernel version 3.2 or higher

- Debian 7 kernel version 3.2 or higher

Note

All of the available file systems for a given Linux distribution may not support hole punching.

How Page Compression Works

When a page is written, it is compressed using the specified compression algorithm. The compressed data is written to disk, where the hole punching mechanism releases empty blocks from the end of the page. If compression fails, data is written out as-is.

Hole Punch Size on Linux

On Linux systems, the file system block size is the unit size used for hole punching. Therefore, page compression only works if page data can be compressed to a size that is less than or equal to the `InnoDB` page size minus the file system block size. For example, if `innodb_page_size=16K` and the file system block size is 4K, page data must compress to less than or equal to 12K to make hole punching possible.

Hole Punch Size on Windows

On Windows systems, the underlying infrastructure for sparse files is based on NTFS compression. Hole punching size is the NTFS compression unit, which is 16 times the NTFS cluster size. Cluster sizes and their compression units are shown in the following table:

Table 15.9 Windows NTFS Cluster Size and Compression Units

Cluster Size	Compression Unit
512 Bytes	8 KB
1 KB	16 KB
2 KB	32 KB
4 KB	64 KB

Page compression on Windows systems only works if page data can be compressed to a size that is less than or equal to the `InnoDB` page size minus the compression unit size.

The default NTFS cluster size is 4K, for which the compression unit size is 64K. This means that page compression has no benefit for an out-of-the box Windows NTFS configuration, as the maximum `innodb_page_size` is also 64K.

For page compression to work on Windows, the file system must be created with a cluster size smaller than 4K, and the `innodb_page_size` must be at least twice the size of the compression unit. For example, for page compression to work on Windows, you could build the file system with a cluster size of 512 Bytes (which has a compression unit of 8KB) and initialize `InnoDB` with an `innodb_page_size` value of 16K or greater.

Enabling Page Compression

To enable page compression, specify the `COMPRESSION` attribute in the `CREATE TABLE` statement. For example:

```
CREATE TABLE t1 (c1 INT) COMPRESSION="zlib";
```

You can also enable page compression in an `ALTER TABLE` statement. However, `ALTER TABLE ... COMPRESSION` only updates the tablespace compression attribute. Writes to the tablespace that occur after setting the new compression algorithm use the new setting, but to apply the new compression algorithm to existing pages, you must rebuild the table using `OPTIMIZE TABLE`.

```
ALTER TABLE t1 COMPRESSION="zlib";
```

```
OPTIMIZE TABLE t1;
```

Disabling Page Compression

To disable page compression, set `COMPRESSION=None` using `ALTER TABLE`. Writes to the tablespace that occur after setting `COMPRESSION=None` no longer use page compression. To uncompress existing pages, you must rebuild the table using `OPTIMIZE TABLE` after setting `COMPRESSION=None`.

```
ALTER TABLE t1 COMPRESSION="None";
OPTIMIZE TABLE t1;
```

Page Compression Metadata

Page compression metadata is found in the `INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES` table, in the following columns:

- `FS_BLOCK_SIZE`: The file system block size, which is the unit size used for hole punching.
- `FILE_SIZE`: The apparent size of the file, which represents the maximum size of the file, uncompressed.
- `ALLOCATED_SIZE`: The actual size of the file, which is the amount of space allocated on disk.

Note

On Unix-like systems, `ls -l` *`tablespace_name.ibd`* shows the apparent file size (equivalent to `FILE_SIZE`) in bytes. To view the actual amount of space allocated on disk (equivalent to `ALLOCATED_SIZE`), use `du --block-size=1` *`tablespace_name.ibd`*. The `--block-size=1` option prints the allocated space in bytes instead of blocks, so that it can be compared to `ls -l` output.

Use `SHOW CREATE TABLE` to view the current page compression setting (`Zlib`, `Lz4`, or `None`). A table may contain a mix of pages with different compression settings.

In the following example, page compression metadata for the employees table is retrieved from the `INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES` table.

```
# Create the employees table with Zlib page compression

CREATE TABLE employees (
    emp_no      INT             NOT NULL,
    birth_date  DATE            NOT NULL,
    first_name  VARCHAR(14)     NOT NULL,
    last_name   VARCHAR(16)     NOT NULL,
    gender      ENUM ('M','F')  NOT NULL,
    hire_date   DATE            NOT NULL,
    PRIMARY KEY (emp_no)
) COMPRESSION="zlib";

# Insert data (not shown)

# Query page compression metadata in INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES

mysql> SELECT SPACE, NAME, FS_BLOCK_SIZE, FILE_SIZE, ALLOCATED_SIZE FROM
INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES WHERE NAME='employees/employees'\G
*************************** 1. row ***************************
SPACE: 45
NAME: employees/employees
FS_BLOCK_SIZE: 4096
FILE_SIZE: 23068672
ALLOCATED_SIZE: 19415040
```

Page compression metadata for the employees table shows that the apparent file size is 23068672 bytes while the actual file size (with page compression) is 19415040 bytes. The file system block size is 4096 bytes, which is the block size used for hole punching.

Page Compression Limitations and Usage Notes

- Page compression is disabled if the file system block size (or compression unit size on Windows) * 2 > `innodb_page_size`.
- Page compression is not supported for tables that reside in shared tablespaces, which include the system tablespace, the temporary tablespace, and general tablespaces.
- Page compression is not supported for undo log tablespaces.
- Page compression is not supported for redo log pages.
- R-tree pages, which are used for spatial indexes, are not compressed.
- Pages that belong to compressed tables (`ROW_FORMAT=COMPRESSED`) are left as-is.
- During recovery, updated pages are written out in an uncompressed form.
- Loading a page-compressed tablespace on a server that does not support the compression algorithm that was used causes an I/O error.
- Before downgrading to an earlier version of MySQL that does not support page compression, uncompress the tables that use the page compression feature. To uncompress a table, run `ALTER TABLE ... COMPRESSION=None` and `OPTIMIZE TABLE`.
- Page-compressed tablespaces can be copied between Linux and Windows servers if the compression algorithm that was used is available on both servers.
- Preserving page compression when moving a page-compressed tablespace file from one host to another requires a utility that preserves sparse files.
- Better page compression may be achieved on Fusion-io hardware with NVMFS than on other platforms, as NVMFS is designed to take advantage of punch hole functionality.
- Using the page compression feature with a large `InnoDB` page size and relatively small file system block size could result in write amplification. For example, a maximum `InnoDB` page size of 64KB with a 4KB file system block size may improve compression but may also increase demand on the buffer pool, leading to increased I/O and potential write amplification.

15.10 InnoDB Row Storage and Row Formats

This section discusses how InnoDB features such as table compression, off-page storage of long variable-length column values, and large index key prefixes are controlled by the row format of an `InnoDB` table. It also discusses considerations for choosing the right row format, and compatibility of row formats between MySQL releases.

15.10.1 Overview of InnoDB Row Storage

The storage for rows and associated columns affects performance for queries and DML operations. As more rows fit into a single disk page, queries and index lookups can work faster, less cache memory is required in the InnoDB buffer pool, and less I/O is required to write out updated values for the numeric and short string columns.

The data in each InnoDB table is divided into pages. The pages that make up each table are arranged in a tree data structure called a B-tree index. Table data and secondary indexes both use this type of structure. The B-tree index that represents an entire table is known as the clustered index, which is organized according to the primary key columns. The nodes of the index data structure contain the values of all the columns in that row (for the clustered index) or the index columns and the primary key columns (for secondary indexes).

Variable-length columns are an exception to this rule. Columns such as `BLOB` and `VARCHAR` that are too long to fit on a B-tree page are stored on separately allocated disk pages called overflow pages.

We call such columns off-page columns. The values of these columns are stored in singly-linked lists of overflow pages, and each such column has its own list of one or more overflow pages. In some cases, all or a prefix of the long column value is stored in the B-tree, to avoid wasting storage and eliminating the need to read a separate page.

The following sections describe how to configure the row format of `InnoDB` tables to control how variable-length columns values are stored. Row format configuration also determines the availability of the table compression feature and large index key prefix support.

15.10.2 Specifying the Row Format for a Table

The default row format is defined by `innodb_default_row_format`, which has a default value of `DYNAMIC`. The default row format is used when the `ROW_FORMAT` table option is not defined explicitly or when `ROW_FORMAT=DEFAULT` is specified.

The row format of a table can be defined explicitly using the `ROW_FORMAT` table option in a `CREATE TABLE` or `ALTER TABLE` statement. For example:

```
CREATE TABLE t1 (c1 INT) ROW_FORMAT=DYNAMIC;
```

An explicitly defined `ROW_FORMAT` setting overrides the implicit default. Specifying `ROW_FORMAT=DEFAULT` is equivalent to using the implicit default.

The `innodb_default_row_format` option can be set dynamically:

```
mysql> SET GLOBAL innodb_default_row_format=DYNAMIC;
```

Valid `innodb_default_row_format` options include `DYNAMIC`, `COMPACT`, and `REDUNDANT`. The `COMPRESSED` row format, which is not supported for use in the system tablespace, cannot be defined as the default. It can only be specified explicitly in a `CREATE TABLE` or `ALTER TABLE` statement. Attempting to set `innodb_default_row_format` to `COMPRESSED` returns an error:

```
mysql> SET GLOBAL innodb_default_row_format=COMPRESSED;
ERROR 1231 (42000): Variable 'innodb_default_row_format'
can't be set to the value of 'COMPRESSED'
```

Newly created tables use the row format defined by `innodb_default_row_format` when a `ROW_FORMAT` option is not specified explicitly or when `ROW_FORMAT=DEFAULT` is used. For example, the following `CREATE TABLE` statements use the row format defined by `innodb_default_row_format`.

```
CREATE TABLE t1 (c1 INT);
```

```
CREATE TABLE t2 (c1 INT) ROW_FORMAT=DEFAULT;
```

When a `ROW_FORMAT` option is not specified explicitly or when `ROW_FORMAT=DEFAULT` is used, any operation that rebuilds a table also silently changes the row format of the table to the format defined by `innodb_default_row_format`.

Table-rebuilding operations include `ALTER TABLE` operations that use `ALGORITHM=COPY` or `ALTER TABLE` operations that use `ALGORITM=INPLACE` where table rebuilding is required. See Table 15.10, “Online Status for DDL Operations” for an overview of the online status of DDL operations. `OPTIMIZE TABLE` is also a table-rebuilding operation.

The following example demonstrates a table-rebuilding operation that silently changes the row format of a table created without an explicitly defined row format.

```
mysql> SELECT @@innodb_default_row_format;
+-----------------------------+
| @@innodb_default_row_format |
+-----------------------------+
| dynamic                     |
+-----------------------------+

mysql> CREATE TABLE t1 (c1 INT);
```

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME LIKE 'test/t1' \G
*************************** 1. row ***************************
     TABLE_ID: 54
         NAME: test/t1
         FLAG: 33
       N_COLS: 4
        SPACE: 35
   ROW_FORMAT: Dynamic
ZIP_PAGE_SIZE: 0
   SPACE_TYPE: Single

mysql> SET GLOBAL innodb_default_row_format=COMPACT;

mysql> ALTER TABLE t1 ADD COLUMN (c2 INT);

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME LIKE 'test/t1' \G
*************************** 1. row ***************************
     TABLE_ID: 55
         NAME: test/t1
         FLAG: 1
       N_COLS: 5
        SPACE: 36
   ROW_FORMAT: Compact
ZIP_PAGE_SIZE: 0
   SPACE_TYPE: Single
```

Consider the following potential issues before changing the row format of existing tables from `REDUNDANT` or `COMPACT` to `DYNAMIC`.

- The `REDUNDANT` and `COMPACT` row format supports a maximum index key prefix length of 767 bytes whereas `DYNAMIC` and `COMPRESSED` row formats support an index key prefix length of 3072 bytes. In a replication environment, if `innodb_default_row_format` is set to `DYNAMIC` on the master and set to `COMPACT` on the slave, the following DDL statement, which does not explicitly define a row format, succeeds on the master but fails on the slave:

  ```
  CREATE TABLE t1 (c1 INT PRIMARY KEY, c2 VARCHAR(5000), KEY i1(c2(3070)));
  ```

  For related information, see Section 15.8.1.7, “Limits on InnoDB Tables”.

- Importing a table that does not explicitly define a row format results in a schema mismatch error if the `innodb_default_row_format` setting on the source server differs from the setting on the destination server. For more information, refer to the limitations outlined in Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance”.

To view the row format of a table, issue a `SHOW TABLE STATUS` statement or query `INFORMATION_SCHEMA.TABLES`.

```
SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME LIKE 'test/t1' \G
```

The row format of an `InnoDB` table determines its physical row structure. See Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table” for more information.

15.10.3 DYNAMIC and COMPRESSED Row Formats

When a table is created with `ROW_FORMAT=DYNAMIC` or `ROW_FORMAT=COMPRESSED`, `InnoDB` can store long variable-length column values (for `VARCHAR`, `VARBINARY`, and `BLOB` and `TEXT` types) fully off-page, with the clustered index record containing only a 20-byte pointer to the overflow page. `InnoDB` also encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.

Whether any columns are stored off-page depends on the page size and the total size of the row. When the row is too long, `InnoDB` chooses the longest columns for off-page storage until the clustered index record fits on the B-tree page. `TEXT` and `BLOB` columns that are less than or equal to 40 bytes are always stored in-line.

The DYNAMIC row format maintains the efficiency of storing the entire row in the index node if it fits (as do the COMPACT and REDUNDANT formats), but the DYNAMIC row format avoids the problem of filling B-tree nodes with a large number of data bytes of long columns. The DYNAMIC format is based on the idea that if a portion of a long data value is stored off-page, it is usually most efficient to store all of the value off-page. With DYNAMIC format, shorter columns are likely to remain in the B-tree node, minimizing the number of overflow pages needed for any given row.

The COMPRESSED row format uses similar internal details for off-page storage as the DYNAMIC row format, with additional storage and performance considerations from the table and index data being compressed and using smaller page sizes. With the COMPRESSED row format, the KEY_BLOCK_SIZE option controls how much column data is stored in the clustered index, and how much is placed on overflow pages. For full details about the COMPRESSED row format, see Section 15.9, “InnoDB Table and Page Compression”.

Both DYNAMIC and COMPRESSED row formats support index key prefixes up to 3072 bytes.

Tables that use the COMPRESSED row format can be created in file-per-table tablespaces or general tablespaces. The system tablespace does not support the COMPRESSED row format. To store a COMPRESSED table in a file-per-table tablespace, innodb_file_per_table must be enabled. The innodb_file_per_table configuration options is not applicable to general tablespaces. General tablespaces support all row formats with the caveat that compressed and uncompressed tables cannot coexist in the same general tablespace due to different physical page sizes. For more information about general tablespaces, see Section 15.7.9, “InnoDB General Tablespaces”.

DYNAMIC tables can be stored in file-per-table tablespaces, general tablespaces, and the system tablespace. To store DYNAMIC tables in the system tablespace, you can either disable innodb_file_per_table and use a regular CREATE TABLE or ALTER TABLE statement, or you can use the TABLESPACE [=] innodb_system table option with CREATE TABLE or ALTER TABLE without having to alter your innodb_file_per_table setting. The innodb_file_per_table configuration option is not applicable to general tablespaces, nor are they applicable when using the TABLESPACE [=] innodb_system table option to store DYNAMIC tables in the system tablespace.

InnoDB does not support compressed temporary tables. When innodb_strict_mode is enabled (the default), CREATE TEMPORARY TABLE returns an error if ROW_FORMAT=COMPRESSED or KEY_BLOCK_SIZE is specified. If innodb_strict_mode is disabled, warnings are issued and the temporary table is created using a non-compressed row format.

DYNAMIC and COMPRESSED row formats are variations of the COMPACT row format and therefore handle CHAR storage in the same way as the COMPACT row format. For more information, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

15.10.4 COMPACT and REDUNDANT Row Formats

InnoDB tables that use the COMPACT or REDUNDANT row format store up to the first 768 bytes of variable-length columns (VARCHAR, VARBINARY, and BLOB and TEXT types) in the index record within the B-tree node, with the remainder stored on the overflow pages. InnoDB also encodes fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored off-page. For example, a CHAR(255) column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with utf8mb4.

For COMPACT or REDUNDANT row formats, if the value of a column is 768 bytes or less, no overflow page is needed, and some savings in I/O may result, since the value is in the B-tree node. This works well for relatively short BLOBs, but may cause B-tree nodes to fill with data rather than key values, reducing their efficiency. Tables with many BLOB columns could cause B-tree nodes to become too full of data, and contain too few rows, making the entire index less efficient than if the rows were shorter or if the column values were stored off-page.

The default row format is DYNAMIC, as defined by the innodb_default_row_format configuration option. See Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats” for more information.

For information about the physical row structure of tables that use the `REDUNDANT` or `COMPACT` row format, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

15.11 InnoDB Disk I/O and File Space Management

As a DBA, you must manage disk I/O to keep the I/O subsystem from becoming saturated, and manage disk space to avoid filling up storage devices. The ACID design model requires a certain amount of I/O that might seem redundant, but helps to ensure data reliability. Within these constraints, `InnoDB` tries to optimize the database work and the organization of disk files to minimize the amount of disk I/O. Sometimes, I/O is postponed until the database is not busy, or until everything needs to be brought to a consistent state, such as during a database restart after a fast shutdown.

This section discusses the main considerations for I/O and disk space with the default kind of MySQL tables (also known as `InnoDB` tables):

- Controlling the amount of background I/O used to improve query performance.
- Enabling or disabling features that provide extra durability at the expense of additional I/O.
- Organizing tables into many small files, a few larger files, or a combination of both.
- Balancing the size of redo log files against the I/O activity that occurs when the log files become full.
- How to reorganize a table for optimal query performance.

15.11.1 InnoDB Disk I/O

`InnoDB` uses asynchronous disk I/O where possible, by creating a number of threads to handle I/O operations, while permitting other database operations to proceed while the I/O is still in progress. On Linux and Windows platforms, `InnoDB` uses the available OS and library functions to perform “native” asynchronous I/O. On other platforms, `InnoDB` still uses I/O threads, but the threads may actually wait for I/O requests to complete; this technique is known as “simulated” asynchronous I/O.

Read-Ahead

If `InnoDB` can determine there is a high probability that data might be needed soon, it performs read-ahead operations to bring that data into the buffer pool so that it is available in memory. Making a few large read requests for contiguous data can be more efficient than making several small, spread-out requests. There are two read-ahead heuristics in `InnoDB`:

- In sequential read-ahead, if `InnoDB` notices that the access pattern to a segment in the tablespace is sequential, it posts in advance a batch of reads of database pages to the I/O system.
- In random read-ahead, if `InnoDB` notices that some area in a tablespace seems to be in the process of being fully read into the buffer pool, it posts the remaining reads to the I/O system.

For information about configuring read-ahead heuristics, see Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”.

Doublewrite Buffer

`InnoDB` uses a novel file flush technique involving a structure called the doublewrite buffer, which is enabled by default in most cases (`innodb_doublewrite=ON`). It adds safety to recovery following a crash or power outage, and improves performance on most varieties of Unix by reducing the need for `fsync()` operations.

Before writing pages to a data file, `InnoDB` first writes them to a contiguous tablespace area called the doublewrite buffer. Only after the write and the flush to the doublewrite buffer has completed does `InnoDB` write the pages to their proper positions in the data file. If there is an operating system, storage subsystem, or `mysqld` process crash in the middle of a page write (causing a torn page condition), `InnoDB` can later find a good copy of the page from the doublewrite buffer during recovery.

If system tablespace files (“ibdata files”) are located on Fusion-io devices that support atomic writes, doublewrite buffering is automatically disabled and Fusion-io atomic writes are used for all data files. Because the doublewrite buffer setting is global, doublewrite buffering is also disabled for data files residing on non-Fusion-io hardware. This feature is only supported on Fusion-io hardware and is only enabled for Fusion-io NVMFS on Linux. To take full advantage of this feature, an `innodb_flush_method` setting of `O_DIRECT` is recommended.

15.11.2 File Space Management

The data files that you define in the configuration file using the `innodb_data_file_path` configuration option form the `InnoDB` system tablespace. The files are logically concatenated to form the system tablespace. There is no striping in use. You cannot define where within the system tablespace your tables are allocated. In a newly created system tablespace, `InnoDB` allocates space starting from the first data file.

To avoid the issues that come with storing all tables and indexes inside the system tablespace, you can enable the `innodb_file_per_table` configuration option (the default), which stores each newly created table in a separate tablespace file (with extension `.ibd`). For tables stored this way, there is less fragmentation within the disk file, and when the table is truncated, the space is returned to the operating system rather than still being reserved by InnoDB within the system tablespace. For more information, see Section 15.7.4, “InnoDB File-Per-Table Tablespaces”.

You can also store tables in general tablespaces. General tablespaces are shared tablespaces created using `CREATE TABLESPACE` syntax. They can be created outside of the MySQL data directory, are capable of holding multiple tables, and support tables of all row formats. For more information, see Section 15.7.9, “InnoDB General Tablespaces”.

Pages, Extents, Segments, and Tablespaces

Each tablespace consists of database pages. Every tablespace in a MySQL instance has the same page size. By default, all tablespaces have a page size of 16KB; you can reduce the page size to 8KB or 4KB by specifying the `innodb_page_size` option when you create the MySQL instance. You can also increase the page size to 32KB or 64KB. For more information, refer to the `innodb_page_size` documentation.

The pages are grouped into extents of size 1MB for pages up to 16KB in size (64 consecutive 16KB pages, or 128 8KB pages, or 256 4KB pages). For a page size of 32KB, extent size is 2MB. For page size of 64KB, extent size is 4MB. The “files” inside a tablespace are called segments in `InnoDB`. (These segments are different from the rollback segment, which actually contains many tablespace segments.)

When a segment grows inside the tablespace, `InnoDB` allocates the first 32 pages to it one at a time. After that, `InnoDB` starts to allocate whole extents to the segment. `InnoDB` can add up to 4 extents at a time to a large segment to ensure good sequentiality of data.

Two segments are allocated for each index in `InnoDB`. One is for nonleaf nodes of the B-tree, the other is for the leaf nodes. Keeping the leaf nodes contiguous on disk enables better sequential I/O operations, because these leaf nodes contain the actual table data.

Some pages in the tablespace contain bitmaps of other pages, and therefore a few extents in an `InnoDB` tablespace cannot be allocated to segments as a whole, but only as individual pages.

When you ask for available free space in the tablespace by issuing a `SHOW TABLE STATUS` statement, `InnoDB` reports the extents that are definitely free in the tablespace. `InnoDB` always reserves some extents for cleanup and other internal purposes; these reserved extents are not included in the free space.

When you delete data from a table, `InnoDB` contracts the corresponding B-tree indexes. Whether the freed space becomes available for other users depends on whether the pattern of deletes frees individual pages or extents to the tablespace. Dropping a table or deleting all rows from it is guaranteed to release the space to other users, but remember that deleted rows are physically removed only

by the purge operation, which happens automatically some time after they are no longer needed for transaction rollbacks or consistent reads. (See Section 15.3, “InnoDB Multi-Versioning”.)

How Pages Relate to Table Rows

The maximum row length is slightly less than half a database page for 4KB, 8KB, 16KB, and 32KB `innodb_page_size` settings. For example, the maximum row length is slightly less than 8KB for the default 16KB `InnoDB` page size. For 64KB pages, the maximum row length is slightly less than 16KB.

If a row does not exceed the maximum row length, all of it is stored locally within the page. If a row exceeds the maximum row length, variable-length columns are chosen for external off-page storage until the row fits within the maximum row length limit. External off-page storage for variable-length columns differs by row format:

- *COMPACT and REDUNDANT Row Formats*

  When a variable-length column is chosen for external off-page storage, `InnoDB` stores the first 768 bytes locally in the row, and the rest externally into overflow pages. Each such column has its own list of overflow pages. The 768-byte prefix is accompanied by a 20-byte value that stores the true length of the column and points into the overflow list where the rest of the value is stored. See Section 15.10.4, “COMPACT and REDUNDANT Row Formats”.

- *DYNAMIC and COMPRESSED Row Formats*

  When a variable-length column is chosen for external off-page storage, `InnoDB` stores a 20-byte pointer locally in the row, and the rest externally into overflow pages. See Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats”.

`LONGBLOB` and `LONGTEXT` columns must be less than 4GB, and the total row length, including `BLOB` and `TEXT` columns, must be less than 4GB.

15.11.3 InnoDB Checkpoints

Making your log files very large may reduce disk I/O during checkpointing. It often makes sense to set the total size of the log files as large as the buffer pool or even larger.

How Checkpoint Processing Works

`InnoDB` implements a checkpoint mechanism known as fuzzy checkpointing. `InnoDB` flushes modified database pages from the buffer pool in small batches. There is no need to flush the buffer pool in one single batch, which would disrupt processing of user SQL statements during the checkpointing process.

During crash recovery, `InnoDB` looks for a checkpoint label written to the log files. It knows that all modifications to the database before the label are present in the disk image of the database. Then `InnoDB` scans the log files forward from the checkpoint, applying the logged modifications to the database.

15.11.4 Defragmenting a Table

Random insertions into or deletions from a secondary index can cause the index to become fragmented. Fragmentation means that the physical ordering of the index pages on the disk is not close to the index ordering of the records on the pages, or that there are many unused pages in the 64-page blocks that were allocated to the index.

One symptom of fragmentation is that a table takes more space than it “should” take. How much that is exactly, is difficult to determine. All `InnoDB` data and indexes are stored in B-trees, and their fill factor may vary from 50% to 100%. Another symptom of fragmentation is that a table scan such as this takes more time than it “should” take:

```
SELECT COUNT(*) FROM t WHERE non_indexed_column <> 12345;
```

The preceding query requires MySQL to perform a full table scan, the slowest type of query for a large table.

To speed up index scans, you can periodically perform a “null” `ALTER TABLE` operation, which causes MySQL to rebuild the table:

```
ALTER TABLE tbl_name ENGINE=INNODB
```

You can also use `ALTER TABLE tbl_name FORCE` to perform a “null” alter operation that rebuilds the table.

Both `ALTER TABLE tbl_name ENGINE=INNODB` and `ALTER TABLE tbl_name FORCE` use online DDL. For more information, see Section 15.12.1, “Online DDL Overview”.

Another way to perform a defragmentation operation is to use `mysqldump` to dump the table to a text file, drop the table, and reload it from the dump file.

If the insertions into an index are always ascending and records are deleted only from the end, the `InnoDB` filespace management algorithm guarantees that fragmentation in the index does not occur.

15.11.5 Reclaiming Disk Space with TRUNCATE TABLE

To reclaim operating system disk space when truncating an `InnoDB` table, the table must be stored in its own .ibd file. For a table to be stored in its own .ibd file, `innodb_file_per_table` must enabled when the table is created. Additionally, there cannot be a foreign key constraint between the table being truncated and other tables, otherwise the `TRUNCATE TABLE` operation fails. A foreign key constraint between two columns in the same table, however, is permitted.

When a table is truncated, it is dropped and re-created in a new `.ibd` file, and the freed space is returned to the operating system. This is in contrast to truncating `InnoDB` tables that are stored within the `InnoDB` system tablespace (tables created when `innodb_file_per_table=OFF`) and tables stored in shared general tablespaces, where only `InnoDB` can use the freed space after the table is truncated.

The ability to truncate tables and return disk space to the operating system also means that physical backups can be smaller. Truncating tables that are stored in the system tablespace (tables created when `innodb_file_per_table=OFF`) or in a general tablespace leaves blocks of unused space in the tablespace.

15.12 InnoDB and Online DDL

The `InnoDB` online DDL feature permits in-place table alterations or concurrent DML, or both. Benefits of this feature include:

- Improved responsiveness and availability in busy production environments, where making a table unavailable for minutes or hours is not practical.
- The ability to adjust the balance between performance and concurrency during a DDL operations using the `LOCK` clause.
  - `LOCK=EXCLUSIVE` blocks access to the table entirely.
  - `LOCK=SHARED` allows queries but not DML.
  - `LOCK=NONE` allows full query and DML access to the table.
  - `LOCK=DEFAULT` or omitting the `LOCK` clause permits as much concurrency as possible depending on the type of DDL operation.
- Avoidance of disk space usage and I/O overhead associated with copying the table and reconstructing secondary indexes.

15.12.1 Online DDL Overview

The online DDL feature enhances many DDL operations that formerly required a table copy or blocked DML operations on the table, or both. Table 15.10, “Online Status for DDL Operations” shows how the online DDL feature applies to each DDL statement.

With the exception of some `ALTER TABLE` partitioning clauses, online DDL operations for partitioned `InnoDB` tables follow the same rules that apply to regular `InnoDB` tables. For more information, see Section 15.12.7, “Online DDL for Partitioned Tables”.

Some factors affect the performance, space usage, and semantics of online DDL operations. For more information, see Section 15.12.8, “Online DDL Limitations”.

- The “In-Place?” column shows which operations permit the `ALGORITHM=INPLACE` clause.
- The “Rebuilds Table?” column shows which operations rebuild the table. For operations that use the `INPLACE` algorithm, the table is rebuilt in place. For operations that do not support the `INPLACE` algorithm, the table copy method is used to rebuild the table.
- The “Permits Concurrent DML?” column shows which operations are performed fully online. You can specify `LOCK=NONE` to assert that concurrent DML is permitted during the DDL operation. MySQL automatically permits concurrent DML when possible.

  Concurrent queries are permitted during all online DDL operations. You can specify `LOCK=SHARED` to assert that concurrent queries are permitted during a DDL operation. MySQL automatically permits concurrent queries when possible.
- The “Notes” column provides additional information and explains exceptions and dependencies related to the “Yes/No” values of other columns. An asterisk indicates an exception or dependency.

Table 15.10 Online Status for DDL Operations

Operation	In-Place?	Rebuilds Table?	Permits Concurrent DML?	Only Modifies Metadata?	Notes
`CREATE INDEX`, `ADD INDEX`	Yes*	No*	Yes	No	Restrictions apply for `FULLTEXT` indexes; see next row.
`ADD FULLTEXT INDEX`	Yes*	No*	No	No	Adding the first `FULLTEXT` index rebuilds the table if there is no user-defined `FTS_DOC_ID` column. Subsequent `FULLTEXT` indexes may be added on the same table without rebuilding the table.
`ADD SPATIAL INDEX`	Yes	No	No	No	
`RENAME INDEX`	Yes	No	Yes	Yes	Only modifies table metadata.
`DROP INDEX`	Yes	No	Yes	Yes	Only modifies table metadata.
`OPTIMIZE TABLE`	Yes*	Yes	Yes	No	In-place operation is not supported for tables with `FULLTEXT` indexes.
Set column default value	Yes	No	Yes	Yes	Only modifies table metadata.
Change auto-increment value	Yes	No	Yes	No*	Modifies a value stored in memory, not the data file.

Operation	In-Place?	Rebuilds Table?	Permits Concurrent DML?	Only Modifies Metadata?	Notes
Add foreign key constraint	Yes*	No	Yes	Yes	The `INPLACE` algorithm is supported when `foreign_key_checks` is disabled. Otherwise, only the `COPY` algorithm is supported.
Drop foreign key constraint	Yes	No	Yes	Yes	`foreign_key_checks` can be enabled or disabled.
Rename column	Yes*	No	Yes*	Yes	To permit concurrent DML, keep the same data type and only change the column name. `ALGORITHM=INPLACE` is not supported for renaming a generated column.
Add column	Yes*	Yes*	Yes*	No	Concurrent DML is not permitted when adding an auto-increment column. Data is reorganized substantially, making it an expensive operation. `ALGORITHM=INPLACE` is supported for adding a virtual generated column but not for adding a stored generated column. Adding a virtual generated column does not require a table rebuild.
Drop column	Yes	Yes*	Yes	No	Data is reorganized substantially, making it an expensive operation. `ALGORITHM=INPLACE` is supported for dropping a generated column. Dropping a virtual generated column does not require a table rebuild.
Reorder columns	Yes	Yes	Yes	No	Data is reorganized substantially, making it an expensive operation.
Change `ROW_FORMAT` property	Yes	Yes	Yes	No	Data is reorganized substantially, making it an expensive operation.
Change `KEY_BLOCK_SIZE` property	Yes	Yes	Yes	No	Data is reorganized substantially, making it an expensive operation.
Make column `NULL`	Yes	Yes*	Yes	No	Rebuilds the table in place. Data is reorganized substantially, making it an expensive operation.
Make column `NOT NULL`	Yes*	Yes*	Yes	No	Rebuilds the table in place. `STRICT_ALL_TABLES` or `STRICT_TRANS_TABLES` `SQL_MODE` is required for the

Operation	In-Place?	Rebuilds Table?	Permits Concurrent DML?	Only Modifies Metadata?	Notes
					operation to succeed. The operation fails if the column contains NULL values. The server prohibits changes to foreign key columns that have the potential to cause loss of referential integrity. See Section 13.1.7, "ALTER TABLE Syntax". Data is reorganized substantially, making it an expensive operation.
Change column data type	No*	Yes	No	No	VARCHAR size may be increased using online ALTER TABLE. See Modifying Column Properties for more information.
Add primary key	Yes*	Yes*	Yes	No	Rebuilds the table in place. Data is reorganized substantially, making it an expensive operation. ALGORITHM=INPLACE is not permitted under certain conditions if columns have to be converted to NOT NULL.
Drop primary key and add another	Yes	Yes	Yes	No	Data is reorganized substantially, making it an expensive operation.
Drop primary key	No	Yes	No	No	Only ALGORITHM=COPY supports dropping a primary key without adding a new one in the same ALTER TABLE statement.
Convert character set	No	Yes*	No	No	Rebuilds the table if the new character encoding is different.
Specify character set	No	Yes*	No	No	Rebuilds the table if the new character encoding is different.
Rebuild with FORCE option	Yes*	Yes	Yes	No	Uses ALGORITHM=INPLACE. ALGORITHM=INPLACE is not supported for tables with FULLTEXT indexes.
"null" rebuild using ALTER TABLE ... ENGINE=INNODB	Yes*	Yes	Yes	No	Uses ALGORITHM=INPLACE. ALGORITHM=INPLACE is not supported for tables with FULLTEXT indexes.
Set STATS_PERSISTENT, STATS_AUTO_RECALC, STATS_SAMPLE_PAGES persistent statistics options	Yes	No	Yes	Yes	Only modifies table metadata.

Operation	In-Place?	Rebuilds Table?	Permits Concurrent DML?	Only Modifies Metadata?	Notes
`ALTER TABLE ... ENCRYPTION`	No	Yes	No	Yes	
Drop a `STORED` column	Yes	Yes*	Yes	No	Rebuilds the table in place.
Modify `STORED` column order	Yes	Yes*	Yes	No	Rebuilds the table in place.
Add a `STORED` column	Yes	Yes*	Yes	No	Rebuilds the table in place.
Drop a `VIRTUAL` column	Yes	No	Yes	Yes	
Modify `VIRTUAL` column order	Yes	No	Yes	Yes	
Add a `VIRTUAL` column	Yes	No	Yes	Yes	

The sections that follow provide basic syntax and usage notes for various online DDL operations.

Adding or Dropping Secondary Indexes

- Adding a secondary index:

```
CREATE INDEX name ON table (col_list);
```

```
ALTER TABLE table ADD INDEX name (col_list);
```

- Dropping a secondary index:

```
DROP INDEX name ON table;
```

```
ALTER TABLE table DROP INDEX name;
```

Although no syntax changes are required in the `CREATE INDEX` or `DROP INDEX` commands, some factors affect the performance, space usage, and semantics of this operation (see Section 15.12.8, “Online DDL Limitations”).

Creating and dropping secondary indexes on `InnoDB` tables skips the table-copying behavior.

The table remains available for read and write operations while the index is being created or dropped. The `CREATE INDEX` or `DROP INDEX` statement only finishes after all transactions that are accessing the table are completed, so that the initial state of the index reflects the most recent contents of the table. Previously, modifying the table while an index is being created or dropped typically resulted in a deadlock that cancelled the `INSERT`, `UPDATE`, or `DELETE` statement on the table.

Online DDL support for adding secondary indexes means that you can generally speed the overall process of creating and loading a table and associated indexes by creating the table without any secondary indexes, then adding the secondary indexes after the data is loaded.

Modifying Column Properties

- Modify the default value for a column:

```
ALTER TABLE tbl ALTER COLUMN col SET DEFAULT literal;
```

```
ALTER TABLE tbl ALTER COLUMN col DROP DEFAULT;
```

The default values for columns are stored in the `InnoDB` data dictionary.

- Changing the auto-increment value for a column:

```
ALTER TABLE table AUTO_INCREMENT=next_value;
```

Especially in a distributed system using replication or sharding, you sometimes reset the auto-increment counter for a table to a specific value. The next row inserted into the table uses the specified value for its auto-increment column. You might also use this technique in a data warehousing environment where you periodically empty all the tables and reload them, and you can restart the auto-increment sequence from 1.

- Renaming a column:

```
ALTER TABLE tbl CHANGE old_col_name new_col_name datatype;
```

When you keep the same data type and `[NOT] NULL` attribute, only changing the column name, this operation can always be performed online.

You can also rename a column that is part of a foreign key constraint. The foreign key definition is automatically updated to use the new column name. Renaming a column participating in a foreign key only works with the in-place mode of `ALTER TABLE`. If you use the `ALGORITHM=COPY` clause, or some other condition causes the command to use `ALGORITHM=COPY` behind the scenes, the `ALTER TABLE` statement fails.

- Extending `VARCHAR` size using an in-place `ALTER TABLE` statement:

```
ALTER TABLE t1 ALGORITHM=INPLACE, CHANGE COLUMN c1 c1 VARCHAR(255);
```

The number of length bytes required by a `VARCHAR` column must remain the same. For `VARCHAR` values of 0 to 255, one length byte is required to encode the value. For `VARCHAR` values of 256 bytes or more, two length bytes are required. As a result, in-place `ALTER TABLE` only supports increasing `VARCHAR` size from 0 to 255 bytes or increasing `VARCHAR` size from a value equal to or greater than 256 bytes. In-place `ALTER TABLE` does not support increasing `VARCHAR` size from less than 256 bytes to a value equal to or greater than 256 bytes. In this case, the number of required length bytes would change from 1 to 2, which is only supported by a table copy (`ALGORITHM=COPY`). For example, attempting to change `VARCHAR` column size from 255 to 256 using in-place `ALTER TABLE` would return an error:

```
ALTER TABLE t1 ALGORITHM=INPLACE, CHANGE COLUMN c1 c1 VARCHAR(256);
ERROR 0A000: ALGORITHM=INPLACE is not supported. Reason: Cannot change
column type INPLACE. Try ALGORITHM=COPY.
```

Decreasing `VARCHAR` size using in-place `ALTER TABLE` is not supported. Decreasing `VARCHAR` size requires a table copy (`ALGORITHM=COPY`).

Adding or Dropping Foreign Keys

- Adding or dropping a foreign key constraint:

```
ALTER TABLE tbl1 ADD CONSTRAINT fk_name FOREIGN KEY index (col1)
  REFERENCES tbl2(col2) referential_actions;
```

```
ALTER TABLE tbl DROP FOREIGN KEY fk_name;
```

Dropping a foreign key can be performed online with the `foreign_key_checks` option enabled or disabled. Creating a foreign key online requires `foreign_key_checks` to be disabled.

If you do not know the names of the foreign key constraints on a particular table, issue the following statement and find the constraint name in the `CONSTRAINT` clause for each foreign key:

```
SHOW CREATE TABLE table\G
```

Or, query the `INFORMATION_SCHEMA.TABLE_CONSTRAINTS` table and use the `CONSTRAINT_NAME` and `CONSTRAINT_TYPE` columns to identify the foreign key names.

You can also drop a foreign key and its associated index in a single statement:

```
ALTER TABLE table DROP FOREIGN KEY constraint, DROP INDEX index;
```

If foreign keys are already present in the table being altered (that is, it is a child table containing a `FOREIGN KEY ... REFERENCE` clause), additional restrictions apply to online DDL operations, even those not directly involving the foreign key columns:

- An `ALTER TABLE` on the child table could wait for another transaction to commit, if a change to the parent table caused associated changes in the child table through an `ON UPDATE` or `ON DELETE` clause using the `CASCADE` or `SET NULL` parameters.
- In the same way, if a table is the parent table in a foreign key relationship, even though it does not contain any `FOREIGN KEY` clauses, it could wait for the `ALTER TABLE` to complete if an `INSERT`, `UPDATE`, or `DELETE` statement caused an `ON UPDATE` or `ON DELETE` action in the child table.

Maintaining CREATE TABLE Statements

As your database schema evolves with new columns, data types, constraints, indexes, and so on, keep your `CREATE TABLE` statements up to date with the latest table definitions. Even with the performance improvements of online DDL, it is more efficient to create stable database structures at the beginning, rather than creating part of the schema and then issuing `ALTER TABLE` statements afterward.

The main exception to this guideline is for secondary indexes on tables with large numbers of rows. It is typically most efficient to create the table with all details specified except the secondary indexes, load the data, then create the secondary indexes. You can use the same technique with foreign keys (load the data first, then set up the foreign keys) if you know the initial data is clean and do not need consistency checks during the loading process.

Whatever sequence of `CREATE TABLE`, `CREATE INDEX`, `ALTER TABLE`, and similar statements went into putting a table together, you can capture the SQL needed to reconstruct the current form of the table by issuing the statement `SHOW CREATE TABLE table\G` (uppercase `\G` required for tidy formatting). This output shows clauses such as numeric precision, `NOT NULL`, and `CHARACTER SET` that are sometimes added behind the scenes, which you may want to leave out when cloning the table on a new system or setting up foreign key columns with identical type.

15.12.2 Online DDL Performance, Concurrency, and Space Requirements

Online DDL improves several aspects of MySQL operation, such as performance, concurrency, availability, and scalability:

- Because queries and DML operations on the table can proceed while the DDL is in progress, applications that access the table are more responsive. Reduced locking and waiting for other resources throughout the MySQL server leads to greater scalability, even for operations not involving the table being altered.
- For in-place operations, by avoiding the disk I/O and CPU cycles to rebuild the table, you minimize the overall load on the database and maintain good performance and high throughput during the DDL operation.
- For in-place operations, because less data is read into the buffer pool than if all the data was copied, you avoid purging frequently accessed data from memory, which formerly could cause a temporary performance dip after a DDL operation.

If an online operation requires temporary sort files, `InnoDB` creates them in the temporary file directory by default, not the directory containing the original table. If this directory is not large enough to hold such files, you may need to set the `tmpdir` system variable to a different directory. Alternatively, you can define a separate temporary directory for `InnoDB` online `ALTER TABLE` operations using the `innodb_tmpdir` configuration option. For more information, see Space Requirements for Online DDL Operations, and Section B.5.3.5, “Where MySQL Stores Temporary Files”.

Locking Options for Online DDL

While an InnoDB table is being changed by a DDL operation, the table may or may not be locked, depending on the internal workings of that operation and the `LOCK` clause of the `ALTER TABLE` statement. By default, MySQL uses as little locking as possible during a DDL operation; you specify the clause either to make the locking more restrictive than it normally would be (thus limiting concurrent DML, or DML and queries), or to ensure that some expected degree of locking is allowed for an operation. If the `LOCK` clause specifies a level of locking that is not available for that specific kind of DDL operation, such as `LOCK=SHARED` or `LOCK=NONE` while creating or dropping a primary key, the clause works like an assertion, causing the statement to fail with an error. The following list shows the different possibilities for the `LOCK` clause, from the most permissive to the most restrictive:

- For DDL operations with `LOCK=NONE`, both queries and concurrent DML are allowed. This clause makes the `ALTER TABLE` fail if the kind of DDL operation cannot be performed with the requested type of locking, so specify `LOCK=NONE` if keeping the table fully available is vital and it is OK to cancel the DDL if that is not possible. For example, you might use this clause in DDLs for tables involving customer signups or purchases, to avoid making those tables unavailable by mistakenly issuing an expensive `ALTER TABLE` statement.

- For DDL operations with `LOCK=SHARED`, any writes to the table (that is, DML operations) are blocked, but the data in the table can be read. This clause makes the `ALTER TABLE` fail if the kind of DDL operation cannot be performed with the requested type of locking, so specify `LOCK=SHARED` if keeping the table available for queries is vital and it is OK to cancel the DDL if that is not possible. For example, you might use this clause in DDLs for tables in a data warehouse, where it is OK to delay data load operations until the DDL is finished, but queries cannot be delayed for long periods.

- For DDL operations with `LOCK=DEFAULT`, or with the `LOCK` clause omitted, MySQL uses the lowest level of locking that is available for that kind of operation, allowing concurrent queries, DML, or both wherever possible. This is the setting to use when making pre-planned, pre-tested changes that you know do not cause any availability problems based on the workload for that table.

- For DDL operations with `LOCK=EXCLUSIVE`, both queries and DML operations are blocked. This clause makes the `ALTER TABLE` fail if the kind of DDL operation cannot be performed with the requested type of locking, so specify `LOCK=EXCLUSIVE` if the primary concern is finishing the DDL in the shortest time possible, and it is OK to make applications wait when they try to access the table. You might also use `LOCK=EXCLUSIVE` if the server is supposed to be idle, to avoid unexpected accesses to the table.

In most cases, an online DDL operation on a table waits for currently executing transactions that are accessing the table to commit or roll back because it requires exclusive access to the table for a brief period while the DDL statement is being prepared. Likewise, the online DDL operation requires exclusive access to the table for a brief time before finishing. Thus, an online DDL statement also waits for transactions that are started while the DDL is in progress to commit or roll back before completing. Consequently, in the case of long running transactions that perform inserts, updates, deletes, or `SELECT ... FOR UPDATE` operations on the table, it is possible for online DDL operation to time out waiting for exclusive access to the table.

A case in which an online DDL operation on a table does not wait for a currently executing transaction to complete can occur when the table is in a foreign key relationship and a transaction is run explicitly on the other table in the foreign key relationship. In this case, the transaction holds an exclusive metadata lock on the table it is updating, but only holds shared `InnoDB` table lock (required for foreign key checking) on the other table. The shared `InnoDB` table lock permits the online DDL operation to proceed but blocks the operation at the commit phase, when an exclusive `InnoDB` table lock is required. This scenario can result in deadlocks as other transactions wait for the online DDL operation to commit. (See Bug #48652, and Bug #77390)

Because there is some processing work involved with recording the changes made by concurrent DML operations, then applying those changes at the end, an online DDL operation could take longer overall than the old-style mechanism that blocks table access from other sessions. The reduction

in raw performance is balanced against better responsiveness for applications that use the table. When evaluating the ideal techniques for changing table structure, consider end-user perception of performance, based on factors such as load times for web pages.

A newly created InnoDB secondary index contains only the committed data in the table at the time the `CREATE INDEX` or `ALTER TABLE` statement finishes executing. It does not contain any uncommitted values, old versions of values, or values marked for deletion but not yet removed from the old index.

Performance of In-Place versus Table-Copying DDL Operations

The raw performance of an online DDL operation is largely determined by whether the operation is performed in-place, or requires copying and rebuilding the entire table. See Table 15.10, “Online Status for DDL Operations” to see what kinds of operations can be performed in-place, and any requirements for avoiding table-copy operations.

The performance speedup from in-place DDL applies to operations on secondary indexes, not to the primary key index. The rows of an InnoDB table are stored in a clustered index organized based on the primary key, forming what some database systems call an “index-organized table”. Because the table structure is closely tied to the primary key, redefining the primary key still requires copying the data.

When an operation on the primary key uses `ALGORITHM=INPLACE`, even though the data is still copied, it is more efficient than using `ALGORITHM=COPY` because:

- No undo logging or associated redo logging is required for `ALGORITHM=INPLACE`. These operations add overhead to DDL statements that use `ALGORITHM=COPY`.
- The secondary index entries are pre-sorted, and so can be loaded in order.
- The change buffer is not used, because there are no random-access inserts into the secondary indexes.

To judge the relative performance of online DDL operations, you can run such operations on a big `InnoDB` table using current and earlier versions of MySQL. You can also run all the performance tests under the latest MySQL version, simulating the previous DDL behavior for the “before” results, by setting the `old_alter_table` system variable. Issue the statement `set old_alter_table=1` in the session, and measure DDL performance to record the “before” figures. Then `set old_alter_table=0` to re-enable the newer, faster behavior, and run the DDL operations again to record the “after” figures.

For a basic idea of whether a DDL operation does its changes in-place or performs a table copy, look at the “rows affected” value displayed after the command finishes. For example, here are lines you might see after doing different types of DDL operations:

- Changing the default value of a column (super-fast, does not affect the table data at all):

```
Query OK, 0 rows affected (0.07 sec)
```

- Adding an index (takes time, but `0 rows affected` shows that the table is not copied):

```
Query OK, 0 rows affected (21.42 sec)
```

- Changing the data type of a column (takes substantial time and does require rebuilding all the rows of the table):

```
Query OK, 1671168 rows affected (1 min 35.54 sec)
```

Note

Changing the data type of a column requires rebuilding all the rows of the table with the exception of changing `VARCHAR` size, which may be performed

using online `ALTER TABLE`. See Modifying Column Properties for more information.

For example, before running a DDL operation on a big table, you might check whether the operation is fast or slow as follows:

1. Clone the table structure.
2. Populate the cloned table with a tiny amount of data.
3. Run the DDL operation on the cloned table.
4. Check whether the “rows affected” value is zero or not. A non-zero value means the operation requires rebuilding the entire table, which might require special planning. For example, you might do the DDL operation during a period of scheduled downtime, or on each replication slave server one at a time.

For a deeper understanding of the reduction in MySQL processing, examine the `performance_schema` and `INFORMATION_SCHEMA` tables related to `InnoDB` before and after DDL operations, to see the number of physical reads, writes, memory allocations, and so on.

Space Requirements for Online DDL Operations

Online DDL operations have the following space requirements:

- Space for temporary log files

  There is one such log file for each index being created or table being altered. This log file stores data inserted, updated, or deleted in the table during the DDL operation. The temporary log file is extended when needed by the value of `innodb_sort_buffer_size`, up to the maximum specified by `innodb_online_alter_log_max_size`. If a temporary log file exceeds the upper size limit, the `ALTER TABLE` operation fails and all uncommitted concurrent DML operations are rolled back. Thus, a large value for this option allows more DML to happen during an online DDL operation, but also extends the period of time at the end of the DDL operation when the table is locked to apply the data from the log.

  If the operation takes so long, and concurrent DML modifies the table so much, that the size of the temporary online log exceeds the value of the `innodb_online_alter_log_max_size` configuration option, the online DDL operation fails with a `DB_ONLINE_LOG_TOO_BIG` error.

- Space for temporary sort files

  Online DDL operations that rebuild the table write temporary sort files to the MySQL temporary directory (`$TMPDIR` on Unix, `%TEMP%` on Windows, or the directory specified by the `--tmpdir` configuration variable) during index creation. Each temporary sort file is large enough to hold all columns defined for the new secondary index plus the columns that are part of the primary key of the clustered index, and each one is removed as soon as it is merged into the final table or index. Such operations may require temporary space equal to the amount of data in the table plus indexes. An online DDL operation that rebuilds the table can cause an error if the operation uses all of the available disk space on the file system where the data directory (`datadir`) resides.

  You can use the `innodb_tmpdir` configuration option to define a separate temporary directory for online DDL operations. The `innodb_tmpdir` option was introduced to help avoid temporary directory overflows that could occur as a result of large temporary sort files created during online `ALTER TABLE` operations that rebuild the table.

- Space for an intermediate table file

  Some online DDL operations that rebuild the table create a temporary intermediate table file in the same directory as the original table as opposed to rebuilding the table in place. An intermediate table file may require space equal to the size of the original table. Operations that rebuild the table in place are noted in Section 15.12.1, “Online DDL Overview”.

15.12.3 Online DDL SQL Syntax

Typically, you do not need to do anything special to enable online DDL when using the `ALTER TABLE` statement for `InnoDB` tables. See Table 15.10, “Online Status for DDL Operations” for the kinds of DDL operations that can be performed in-place, allowing concurrent DML, or both. Some variations require particular combinations of configuration settings or `ALTER TABLE` clauses.

You can control the various aspects of a particular online DDL operation by using the `LOCK` and `ALGORITHM` clauses of the `ALTER TABLE` statement. These clauses come at the end of the statement, separated from the table and column specifications by commas. The `LOCK` clause is useful for fine-tuning the degree of concurrent access to the table. The `ALGORITHM` clause is primarily intended for performance comparisons and as a fallback to the older table-copying behavior in case you encounter any issues with existing DDL code. For example:

- To avoid accidentally making the table unavailable for reads, writes, or both, specify a clause on the `ALTER TABLE` statement such as `LOCK=NONE` (allow both reads and writes) or `LOCK=SHARED` (allow reads). The operation halts immediately if the requested level of concurrency is not available.
- To compare performance, run one statement with `ALGORITHM=INPLACE` and another with `ALGORITHM=COPY`, as an alternative to setting the `old_alter_table` configuration option.
- To avoid tying up the server with an `ALTER TABLE` operation that copies the table, include `ALGORITHM=INPLACE`. The statement halts immediately if it cannot use the in-place mechanism. See Table 15.10, “Online Status for DDL Operations” for a list of the DDL operations that can or cannot be performed in-place.

See Section 15.12.2, “Online DDL Performance, Concurrency, and Space Requirements” for more details about the `LOCK` clause.

15.12.4 Simplifying DDL Statements with Online DDL

Before the introduction of online DDL, it was common practice to combine many DDL operations into a single `ALTER TABLE` statement. Because each `ALTER TABLE` statement involved copying and rebuilding the table, it was more efficient to make several changes to the same table at once, since those changes could all be done with a single rebuild operation for the table. The downside was that SQL code involving DDL operations was harder to maintain and to reuse in different scripts. If the specific changes were different each time, you might have to construct a new complex `ALTER TABLE` for each slightly different scenario.

For DDL operations that can be done in-place, as shown in Table 15.10, “Online Status for DDL Operations”, now you can separate them into individual `ALTER TABLE` statements for easier scripting and maintenance, without sacrificing efficiency. For example, you might take a complicated statement such as:

```
ALTER TABLE t1 ADD INDEX i1(c1), ADD UNIQUE INDEX i2(c2),
  CHANGE c4_old_name c4_new_name INTEGER UNSIGNED;
```

and break it down into simpler parts that can be tested and performed independently, such as:

```
ALTER TABLE t1 ADD INDEX i1(c1);
ALTER TABLE t1 ADD UNIQUE INDEX i2(c2);
ALTER TABLE t1 CHANGE c4_old_name c4_new_name INTEGER UNSIGNED NOT NULL;
```

You might still use multi-part `ALTER TABLE` statements for:

- Operations that must be performed in a specific sequence, such as creating an index followed by a foreign key constraint that uses that index.
- Operations all using the same specific `LOCK` clause, that you want to either succeed or fail as a group.

- Operations that cannot be performed in-place, that is, that still copy and rebuild the table.
- Operations for which you specify `ALGORITHM=COPY` or `old_alter_table=1`, to force the table-copying behavior if needed for precise backward-compatibility in specialized scenarios.

15.12.5 Online DDL Implementation Details

Each `ALTER TABLE` operation for an `InnoDB` table is governed by several aspects:

- Whether there is any change to the physical representation of the table, or whether it purely a change to metadata that can be done without touching the table itself.
- Whether the volume of data in the table stays the same, increases, or decreases.
- Whether a change in table data involves the clustered index, secondary indexes, or both.
- Whether there are any foreign key relationships between the table being altered and some other table. The mechanics differ depending on whether the `foreign_key_checks` configuration option is enabled or disabled.
- Whether the table is partitioned. Partitioning clauses of `ALTER TABLE` are turned into low-level operations involving one or more tables, and those operations follow the regular rules for online DDL.
- Whether the table data must be copied, whether the table can be reorganized “in-place”, or a combination of both.
- Whether the table contains any auto-increment columns.
- What degree of locking is required, either by the nature of the underlying database operations, or a `LOCK` clause that you specify in the `ALTER TABLE` statement.

Error Conditions for Online DDL

Here are the primary reasons why an online DDL operation could fail:

- If a `LOCK` clause specifies a low degree of locking (`SHARED` or `NONE`) that is not compatible with the particular type of DDL operation.
- If a timeout occurs while waiting to get an exclusive lock on the table, which may be needed briefly during the initial and final phases of the DDL operation.
- If the `tmpdir` or `innodb_tmpdir` file system runs out of disk space, while MySQL writes temporary sort files on disk during index creation. For more information, see Section B.5.3.5, “Where MySQL Stores Temporary Files”.
- If the `ALTER TABLE` takes so long, and concurrent DML modifies the table so much, that the size of the temporary online log exceeds the value of the `innodb_online_alter_log_max_size` configuration option. This condition causes a `DB_ONLINE_LOG_TOO_BIG` error.
- If concurrent DML makes changes to the table that are allowed with the original table definition, but not with the new one. The operation only fails at the very end, when MySQL tries to apply all the changes from concurrent DML statements. For example, you might insert duplicate values into a column while a unique index is being created, or you might insert `NULL` values into a column while creating a primary key index on that column. The changes made by the concurrent DML take precedence, and the `ALTER TABLE` operation is effectively rolled back.

Although the configuration option `innodb_file_per_table` has a dramatic effect on the representation for an `InnoDB` table, all online DDL operations work equally well whether that option is enabled or disabled, and whether the table is physically located in its own .ibd file or inside the system tablespace.

InnoDB has two types of indexes: the clustered index representing all the data in the table, and optional secondary indexes to speed up queries. Since the clustered index contains the data values in its B-tree

nodes, adding or dropping a clustered index does involve copying the data, and creating a new copy of the table. A secondary index, however, contains only the index key and the value of the primary key. This type of index can be created or dropped without copying the data in the clustered index. Because each secondary index contains copies of the primary key values (used to access the clustered index when needed), when you change the definition of the primary key, all secondary indexes are recreated as well.

Dropping a secondary index is simple. Only the internal InnoDB system tables and the MySQL data dictionary tables are updated to reflect the fact that the index no longer exists. InnoDB returns the storage used for the index to the tablespace that contained it, so that new indexes or additional table rows can use the space.

To add a secondary index to an existing table, InnoDB scans the table, and sorts the rows using memory buffers and temporary files in order by the values of the secondary index key columns. The B-tree is then built in key-value order, which is more efficient than inserting rows into an index in random order. Because the B-tree nodes are split when they fill, building the index in this way results in a higher fill-factor for the index, making it more efficient for subsequent access.

15.12.6 Online DDL and Crash Recovery

Although no data is lost if the server crashes while an `ALTER TABLE` statement is executing, the crash recovery process is different for clustered indexes and secondary indexes.

If the server crashes while creating an InnoDB secondary index, upon recovery, MySQL drops any partially created indexes. You must re-run the `ALTER TABLE` or `CREATE INDEX` statement.

When a crash occurs during the creation of an InnoDB clustered index, recovery is more complicated, because the data in the table must be copied to an entirely new clustered index. Remember that all InnoDB tables are stored as clustered indexes.

MySQL creates the new clustered index by copying the existing data from the original InnoDB table to a temporary table that has the desired index structure. Once the data is completely copied to this temporary table, the original table is renamed with a different temporary table name. The temporary table comprising the new clustered index is renamed with the name of the original table, and the original table is dropped from the database.

If a system crash occurs while creating a new clustered index, no data is lost, but you must complete the recovery process using the temporary tables that exist during the process. Since it is rare to re-create a clustered index or re-define primary keys on large tables, or to encounter a system crash during this operation, this manual does not provide information on recovering from this scenario.

15.12.7 Online DDL for Partitioned Tables

Some `ALTER TABLE` partitioning clauses do not go through the same internal online DDL API as regular non-partitioned `InnoDB` tables. As a result, online support for `ALTER TABLE` partitioning clauses varies.

The following table shows the online status for each `ALTER TABLE` partitioning statement. Regardless of the online DDL API that is used, MySQL attempts to minimize data copying and locking where possible.

- The “In-Place?” column shows which operations permit the `ALGORITHM=INPLACE` clause.
- The “Permits Concurrent DML?” column shows which operations are performed fully online. You can specify `LOCK=NONE` to assert that concurrent DML is permitted during the DDL operation. MySQL automatically permits concurrent DML where possible.

  For operations that support `ALGORITHM={COPY|INPLACE}`, you can specify `LOCK=SHARED` to assert that concurrent queries are permitted during a DDL operation. MySQL automatically permits concurrent queries where possible.

- The “Notes” column provides additional information and explains exceptions and dependencies related to the “Yes/No” values of other columns. An asterisk indicates an exception or dependency.

`ALTER TABLE` partitioning options that use `ALGORITHM=COPY` or that only permit “`ALGORITHM=DEFAULT, LOCK=DEFAULT`”, repartition the table using the `COPY` algorithm. In other words, a new partitioned table is created with the new partitioning scheme. The newly created table includes any changes applied by the `ALTER TABLE` statement, and table data is copied into the new table structure.

Table 15.11 Online Status for ALTER TABLE Partitioning Clauses

Partitioning Clause	In-Place?	Permits Concurrent DML?	Notes
`PARTITION BY`	No	No	Permits `ALGORITHM=COPY`, `LOCK={DEFAULT|SHARED|EXCLUSIVE}`
`ADD PARTITION`	Yes*	Yes*	Prior to 8.0.2, only `ALGORITHM=DEFAULT, LOCK=DEFAULT` is supported. As of 8.0.2, `ALGORITHM=INPLACE, LOCK={DEFAULT|NONE|SHARED|EXCLUSISVE}` is supported for `RANGE` and `LIST` partitions, `ALGORITHM=INPLACE, LOCK={DEFAULT|SHARED|EXCLUSISVE}` for `HASH` and `KEY` partitions, and `ALGORITHM=COPY, LOCK={SHARED|EXCLUSIVE}` for all partition types. Does not copy existing data for tables partitioned by `RANGE` or `LIST`. Concurrent queries are permitted with `ALGORITHM=COPY` for tables partitioned by `HASH` or `LIST`, as MySQL copies the data while holding a shared lock.
`DROP PARTITION`	Yes*	Yes*	Prior to 8.0.2, only `ALGORITHM=DEFAULT, LOCK=DEFAULT` is supported. As of 8.0.2, `ALGORITHM=INPLACE, LOCK={DEFAULT|NONE|SHARED|EXCLUSIVE}` is supported. Does not copy data for tables partitioned by `RANGE` or `LIST`. `DROP PARTITION` with `ALGORITHM=INPLACE` deletes data stored in the partition and drops the partition. However, `DROP PARTITION` with `ALGORITHM=COPY` or `old_alter_table=ON` rebuilds the partitioned table and attempts to move data from the dropped partition to another partition with a compatible `PARTITION ... VALUES` definition. Data that cannot be moved to another partition is deleted.
`DISCARD PARTITION`	No	No	Only permits `ALGORITHM=DEFAULT`, `LOCK=DEFAULT`

Partitioning Clause	In-Place?	Permits Concurrent DML?	Notes
`IMPORT PARTITION`	No	No	Only permits `ALGORITHM=DEFAULT`, `LOCK=DEFAULT`
`TRUNCATE PARTITION`	Yes	Yes	Does not copy existing data. It merely deletes rows; it does not alter the definition of the table itself, or of any of its partitions.
`COALESCE PARTITION`	Yes*	No	Prior to 8.0.2, only `ALGORITHM=DEFAULT, LOCK=DEFAULT` is supported, and concurrent queries are permitted for tables partitioned by `HASH` or `LIST`. MySQL copies the data while holding a shared lock. As of 8.0.2, `ALGORITHM=INPLACE, LOCK={DEFAULT|SHARED|EXCLUSIVE}` is supported.
`REORGANIZE PARTITION`	Yes*	No	Prior to 8.0.2, only `ALGORITHM=DEFAULT, LOCK=DEFAULT` is supported, and concurrent queries are permitted for tables partitioned by `LINEAR HASH` or `LIST`. MySQL copies data from affected partitions while holding a shared metadata lock. As of 8.0.2, `ALGORITHM=INPLACE, LOCK={DEFAULT|SHARED|EXCLUSIVE}` is supported.
`EXCHANGE PARTITION`	Yes	Yes	
`ANALYZE PARTITION`	Yes	Yes	
`CHECK PARTITION`	Yes	Yes	
`OPTIMIZE PARTITION`	No	No	`ALGORITHM` and `LOCK` clauses are ignored. Rebuilds the entire table. See Section 22.3.4, “Maintenance of Partitions”.
`REBUILD PARTITION`	Yes*	No	Prior to 8.0.2, only `ALGORITHM=DEFAULT, LOCK=DEFAULT` is supported, and concurrent queries are permitted for tables partitioned by `LINEAR HASH` or `LIST`. MySQL copies data from affected partitions while holding a shared metadata lock. As of 8.0.2, `ALGORITHM=INPLACE, LOCK={DEFAULT|SHARED|EXCLUSIVE}` is supported.
`REPAIR PARTITION`	Yes	Yes	
`REMOVE PARTITIONING`	No	No	Permits `ALGORITHM=COPY`, `LOCK={DEFAULT|SHARED|EXCLUSIVE}`

Non-partitioning online `ALTER TABLE` operations on partitioned tables follow the same rules that apply to regular tables. However, `ALTER TABLE` performs online operations on each table partition, which causes increased demand on system resources due to operations being performed on multiple partitions.

For additional information about `ALTER TABLE` partitioning clauses, see Partitioning Options, and Section 13.1.7.1, “ALTER TABLE Partition Operations”. For information about partitioning in general, see Chapter 22, *Partitioning*.

15.12.8 Online DDL Limitations

The following limitations apply to online DDL operations:

- The table is copied, rather than using Fast Index Creation when you create an index on a `TEMPORARY TABLE`. This has been reported as MySQL Bug #39833.
- InnoDB handles error cases when users attempt to drop indexes needed for foreign keys. See Section B.3, “Server Error Codes and Messages” for information related to error `1553`.
- The `ALTER TABLE` clause `LOCK=NONE` is not allowed if there are `ON...CASCADE` or `ON...SET NULL` constraints on the table.
- Depending on the internal workings of the online DDL operation and the `LOCK` clause of the `ALTER TABLE` statement, an online DDL operation may require exclusive access to the table for a brief time during the initial and final phases of the DDL operation. Thus, an online DDL operation might wait before finishing if there is a long-running transaction performing inserts, updates, deletes, or `SELECT ... FOR UPDATE` on the table; and an online DDL operation might wait before finishing if a similar long-running transaction is started while the `ALTER TABLE` is in progress.
- When running an online DDL operation, the thread that runs the `ALTER TABLE` statement applies an “online log” of DML operations that were run concurrently on the same table from other connection threads. When the DML operations are applied, it is possible to encounter a duplicate key entry error (`ERROR 1062 (23000): Duplicate entry`), even if the duplicate entry is only temporary and would be reverted by a later entry in the “online log”. This is similar to the idea of a foreign key constraint check in `InnoDB` in which constraints must hold during a transaction.
- `OPTIMIZE TABLE` for an `InnoDB` table is mapped to an `ALTER TABLE` operation to rebuild the table and update index statistics and free unused space in the clustered index. Secondary indexes are not created as efficiently because keys are inserted in the order they appeared in the primary key. `OPTIMIZE TABLE` also supports online DDL for rebuilding regular and partitioned `InnoDB` tables. For additional information, see Section 15.12.1, “Online DDL Overview”.

  Note

  Prior to MySQL 5.6.17 / 5.7.4, there was not online DDL support for this operation.

- `InnoDB` tables created before MySQL 5.6 do not support `ALTER TABLE ... ALGORITHM=INPLACE` for tables that include temporal columns (`DATE`, `DATETIME` or `TIMESTAMP`) and have not been rebuilt using `ALTER TABLE ... ALGORITHM=COPY`. In this case, an `ALTER TABLE ... ALGORITHM=INPLACE` operation returns the following error:

  ```
  ERROR 1846 (0A000): ALGORITHM=INPLACE is not supported.
  Reason: Cannot change column type INPLACE. Try ALGORITHM=COPY.
  ```

- These limitations are generally applicable to online DDL operations on large tables where table copying is involved:
  - There is no mechanism to pause an online DDL operation or to throttle I/O or CPU usage for an online DDL operation.
  - Rollback of an online DDL operation can be expensive should the operation fail.
  - Long running online DDL operations can cause replication lag. An online DDL operation must finish running on the master before it is run on the slave. Also, DML that was processed

concurrently on the master is only processed on the slave after the DDL operation on the slave is completed (Bug #73196).

For additional information related to running online DDL operations on large tables, see Section 15.12.2, “Online DDL Performance, Concurrency, and Space Requirements”.

15.13 InnoDB Startup Options and System Variables

- System variables that are true or false can be enabled at server startup by naming them, or disabled by using a `--skip-` prefix. For example, to enable or disable the `InnoDB` adaptive hash index, you can use `--innodb_adaptive_hash_index` or `--skip-innodb_adaptive_hash_index` on the command line, or `innodb_adaptive_hash_index` or `skip-innodb_adaptive_hash_index` in an option file.
- System variables that take a numeric value can be specified as `--var_name=value` on the command line or as `var_name=value` in option files.
- Many system variables can be changed at runtime (see Section 5.1.6.2, “Dynamic System Variables”).
- For information about `GLOBAL` and `SESSION` variable scope modifiers, refer to the `SET` statement documentation.
- Certain options control the locations and layout of the `InnoDB` data files. Section 15.6.1, “InnoDB Startup Configuration” explains how to use these options.
- Some options, which you might not use initially, help tune `InnoDB` performance characteristics based on machine capacity and your database workload.
- For more information on specifying options and system variables, see Section 4.2.3, “Specifying Program Options”.

Table 15.12 `InnoDB` Option/Variable Reference

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
daemon_memcached_enable_binlog	Yes	Yes	Yes		Global	No
daemon_memcached_engine_lib_name	Yes	Yes	Yes		Global	No
daemon_memcached_engine_lib_path	Yes	Yes	Yes		Global	No
daemon_memcached_option	Yes	Yes	Yes		Global	No
daemon_memcached_r_batch_size	Yes	Yes	Yes		Global	No
daemon_memcached_w_batch_size	Yes	Yes	Yes		Global	No
foreign_key_checks			Yes		Both	Yes
ignore-builtin-innodb	Yes	Yes			Global	No
- *Variable*: ignore_builtin_innodb			Yes		Global	No
innodb	Yes	Yes				
innodb_adaptive_flushing	Yes	Yes	Yes		Global	Yes
innodb_adaptive_flushing_lwm	Yes	Yes	Yes		Global	Yes
innodb_adaptive_hash_index	Yes	Yes	Yes		Global	Yes
innodb_adaptive_hash_index_parts	Yes	Yes	Yes		Global	No
innodb_adaptive_max_sleep_delay	Yes	Yes	Yes		Global	Yes
innodb_api_bk_commit_interval	Yes	Yes	Yes		Global	Yes
innodb_api_disable_rowlock	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_api_enable_binlog	Yes	Yes	Yes		Global	No
innodb_api_enable_mdl	Yes	Yes	Yes		Global	No
innodb_api_trx_level	Yes	Yes	Yes		Global	Yes
innodb_autoextend_increment	Yes	Yes	Yes		Global	Yes
innodb_autoinc_lock_mode	Yes	Yes	Yes		Global	No
Innodb_available_undo_logs				Yes	Global	No
innodb_background_drop_list_empty	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_bytes_data				Yes	Global	No
Innodb_buffer_pool_bytes_dirty				Yes	Global	No
innodb_buffer_pool_chunk_size	Yes	Yes	Yes		Global	No
innodb_buffer_pool_debug	Yes	Yes	Yes		Global	No
innodb_buffer_pool_dump_at_shutdown	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_dump_now	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_dump_pct	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_dump_status				Yes	Global	No
innodb_buffer_pool_filename	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_instances	Yes	Yes	Yes		Global	No
innodb_buffer_pool_load_abort	Yes	Yes	Yes		Global	Yes
innodb_buffer_pool_load_at_startup	Yes	Yes	Yes		Global	No
innodb_buffer_pool_load_now	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_load_status				Yes	Global	No
Innodb_buffer_pool_pages_data				Yes	Global	No
Innodb_buffer_pool_pages_dirty				Yes	Global	No
Innodb_buffer_pool_pages_flushed				Yes	Global	No
Innodb_buffer_pool_pages_free				Yes	Global	No
Innodb_buffer_pool_pages_latched				Yes	Global	No
Innodb_buffer_pool_pages_misc				Yes	Global	No
Innodb_buffer_pool_pages_total				Yes	Global	No
Innodb_buffer_pool_read_ahead				Yes	Global	No
Innodb_buffer_pool_read_ahead_evicted				Yes	Global	No
Innodb_buffer_pool_read_ahead_rnd				Yes	Global	No
Innodb_buffer_pool_read_requests				Yes	Global	No
Innodb_buffer_pool_reads				Yes	Global	No
Innodb_buffer_pool_resize_status				Yes	Global	No
innodb_buffer_pool_size	Yes	Yes	Yes		Global	Yes
Innodb_buffer_pool_wait_free				Yes	Global	No
Innodb_buffer_pool_write_requests				Yes	Global	No
innodb_change_buffer_max_size	Yes	Yes	Yes		Global	Yes
innodb_change_buffering	Yes	Yes	Yes		Global	Yes
innodb_change_buffering_debug	Yes	Yes	Yes		Global	Yes
innodb_checkpoint_disabled	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_checksum_algorithm	Yes	Yes	Yes		Global	Yes
innodb_cmp_per_index_enabled	Yes	Yes	Yes		Global	Yes
innodb_commit_concurrency	Yes	Yes	Yes		Global	Yes
innodb_compress_debug	Yes	Yes	Yes		Global	Yes
innodb_compression_failure_threshold_pct	Yes	Yes	Yes		Global	Yes
innodb_compression_level	Yes	Yes	Yes		Global	Yes
innodb_compression_pad_pct_max	Yes	Yes	Yes		Global	Yes
innodb_concurrency_tickets	Yes	Yes	Yes		Global	Yes
innodb_data_file_path	Yes	Yes	Yes		Global	No
Innodb_data_fsyncs				Yes	Global	No
innodb_data_home_dir	Yes	Yes	Yes		Global	No
Innodb_data_pending_fsyncs				Yes	Global	No
Innodb_data_pending_reads				Yes	Global	No
Innodb_data_pending_writes				Yes	Global	No
Innodb_data_read				Yes	Global	No
Innodb_data_reads				Yes	Global	No
Innodb_data_writes				Yes	Global	No
Innodb_data_written				Yes	Global	No
Innodb_dblwr_pages_written				Yes	Global	No
Innodb_dblwr_writes				Yes	Global	No
innodb_default_row_format	Yes	Yes	Yes		Global	Yes
innodb_disable_sort_file_cache	Yes	Yes	Yes		Global	Yes
innodb_doublewrite	Yes	Yes	Yes		Global	No
innodb_fast_shutdown	Yes	Yes	Yes		Global	Yes
innodb_fil_make_page_dirty_debug	Yes	Yes	Yes		Global	Yes
innodb_file_per_table	Yes	Yes	Yes		Global	Yes
innodb_fill_factor	Yes	Yes	Yes		Global	Yes
innodb_flush_log_at_timeout			Yes		Global	Yes
innodb_flush_log_at_trx_commit	Yes	Yes	Yes		Global	Yes
innodb_flush_method	Yes	Yes	Yes		Global	No
innodb_flush_neighbors	Yes	Yes	Yes		Global	Yes
innodb_flush_sync	Yes	Yes	Yes		Global	Yes
innodb_flushing_avg_loops	Yes	Yes	Yes		Global	Yes
innodb_force_load_corrupted	Yes	Yes	Yes		Global	No
innodb_force_recovery	Yes	Yes	Yes		Global	No
innodb_ft_aux_table	Yes	Yes	Yes		Global	Yes
innodb_ft_cache_size	Yes	Yes	Yes		Global	No
innodb_ft_enable_diag_print	Yes	Yes	Yes		Global	Yes
innodb_ft_enable_stopword	Yes	Yes	Yes		Global	Yes
innodb_ft_max_token_size	Yes	Yes	Yes		Global	No
innodb_ft_min_token_size	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_ft_num_word_optimize	Yes	Yes	Yes		Global	Yes
innodb_ft_result_cache_limit	Yes	Yes	Yes		Global	Yes
innodb_ft_server_stopword_table	Yes	Yes	Yes		Global	Yes
innodb_ft_sort_pll_degree	Yes	Yes	Yes		Global	No
innodb_ft_total_cache_size	Yes	Yes	Yes		Global	No
innodb_ft_user_stopword_table	Yes	Yes	Yes		Both	Yes
Innodb_have_atomic_builtins				Yes	Global	No
innodb_io_capacity	Yes	Yes	Yes		Global	Yes
innodb_io_capacity_max	Yes	Yes	Yes		Global	Yes
innodb_limit_optimistic_insert_debug	Yes	Yes	Yes		Global	Yes
innodb_lock_wait_timeout	Yes	Yes	Yes		Both	Yes
innodb_log_buffer_size	Yes	Yes	Yes		Global	No
innodb_log_checksums	Yes	Yes	Yes		Global	Yes
innodb_log_compressed_pages	Yes	Yes	Yes		Global	Yes
innodb_log_file_size	Yes	Yes	Yes		Global	No
innodb_log_files_in_group	Yes	Yes	Yes		Global	No
innodb_log_group_home_dir	Yes	Yes	Yes		Global	No
Innodb_log_waits				Yes	Global	No
innodb_log_write_ahead_size	Yes	Yes	Yes		Global	Yes
Innodb_log_write_requests				Yes	Global	No
Innodb_log_writes				Yes	Global	No
innodb_lru_scan_depth	Yes	Yes	Yes		Global	Yes
innodb_max_dirty_pages_pct	Yes	Yes	Yes		Global	Yes
innodb_max_dirty_pages_pct_lwm	Yes	Yes	Yes		Global	Yes
innodb_max_purge_lag	Yes	Yes	Yes		Global	Yes
innodb_max_purge_lag_delay	Yes	Yes	Yes		Global	Yes
innodb_max_undo_log_size	Yes	Yes	Yes		Global	Yes
innodb_merge_threshold_set_all_debug	Yes	Yes	Yes		Global	Yes
innodb_monitor_disable	Yes	Yes	Yes		Global	Yes
innodb_monitor_enable	Yes	Yes	Yes		Global	Yes
innodb_monitor_reset	Yes	Yes	Yes		Global	Yes
innodb_monitor_reset_all	Yes	Yes	Yes		Global	Yes
Innodb_num_open_files				Yes	Global	No
innodb_numa_interleave	Yes	Yes	Yes		Global	No
innodb_old_blocks_pct	Yes	Yes	Yes		Global	Yes
innodb_old_blocks_time	Yes	Yes	Yes		Global	Yes
innodb_online_alter_log_max_size	Yes	Yes	Yes		Global	Yes
innodb_open_files	Yes	Yes	Yes		Global	No
innodb_optimize_fulltext_only	Yes	Yes	Yes		Global	Yes
Innodb_os_log_fsyncs				Yes	Global	No
Innodb_os_log_pending_fsyncs				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Innodb_os_log_pending_writes				Yes	Global	No
Innodb_os_log_written				Yes	Global	No
innodb_page_cleaners	Yes	Yes	Yes		Global	No
Innodb_page_size				Yes	Global	No
innodb_page_size	Yes	Yes	Yes		Global	No
Innodb_pages_created				Yes	Global	No
Innodb_pages_read				Yes	Global	No
Innodb_pages_written				Yes	Global	No
innodb_print_all_deadlocks	Yes	Yes	Yes		Global	Yes
innodb_purge_batch_size	Yes	Yes	Yes		Global	Yes
innodb_purge_rseg_truncate_frequency	Yes	Yes	Yes		Global	Yes
innodb_purge_threads	Yes	Yes	Yes		Global	No
innodb_random_read_ahead	Yes	Yes	Yes		Global	Yes
innodb_read_ahead_threshold	Yes	Yes	Yes		Global	Yes
innodb_read_io_threads	Yes	Yes	Yes		Global	No
innodb_read_only	Yes	Yes	Yes		Global	No
innodb_redo_log_encrypt	Yes	Yes	Yes		Global	Yes
innodb_replication_delay	Yes	Yes	Yes		Global	Yes
innodb_rollback_on_timeout	Yes	Yes	Yes		Global	No
innodb_rollback_segments	Yes	Yes	Yes		Global	Yes
Innodb_row_lock_current_waits				Yes	Global	No
Innodb_row_lock_time				Yes	Global	No
Innodb_row_lock_time_avg				Yes	Global	No
Innodb_row_lock_time_max				Yes	Global	No
Innodb_row_lock_waits				Yes	Global	No
Innodb_rows_deleted				Yes	Global	No
Innodb_rows_inserted				Yes	Global	No
Innodb_rows_read				Yes	Global	No
Innodb_rows_updated				Yes	Global	No
innodb_saved_page_number_debug	Yes	Yes	Yes		Global	Yes
innodb_scan_directories	Yes	Yes	Yes		Global	No
innodb_sort_buffer_size	Yes	Yes	Yes		Global	No
innodb_spin_wait_delay	Yes	Yes	Yes		Global	Yes
innodb_stats_auto_recalc	Yes	Yes	Yes		Global	Yes
innodb_stats_include_delete_marked	Yes	Yes	Yes		Global	Yes
innodb_stats_method	Yes	Yes	Yes		Global	Yes
innodb_stats_on_metadata	Yes	Yes	Yes		Global	Yes
innodb_stats_persistent	Yes	Yes	Yes		Global	Yes
innodb_stats_persistent_sample_pages	Yes	Yes	Yes		Global	Yes
innodb_stats_transient_sample_pages	Yes	Yes	Yes		Global	Yes
innodb-status-file	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
innodb_status_output	Yes	Yes	Yes		Global	Yes
innodb_status_output_locks	Yes	Yes	Yes		Global	Yes
innodb_strict_mode	Yes	Yes	Yes		Both	Yes
innodb_sync_array_size	Yes	Yes	Yes		Global	No
innodb_sync_debug	Yes	Yes	Yes		Global	No
innodb_sync_spin_loops	Yes	Yes	Yes		Global	Yes
innodb_table_locks	Yes	Yes	Yes		Both	Yes
innodb_temp_data_file_path	Yes	Yes	Yes		Global	No
innodb_thread_concurrency	Yes	Yes	Yes		Global	Yes
innodb_thread_sleep_delay	Yes	Yes	Yes		Global	Yes
innodb_tmpdir	Yes	Yes	Yes		Both	Yes
Innodb_truncated_status_writes				Yes	Global	No
innodb_trx_purge_view_update_only_debug	Yes	Yes	Yes		Global	Yes
innodb_trx_rseg_n_slots_debug	Yes	Yes	Yes		Global	Yes
innodb_undo_directory	Yes	Yes	Yes		Global	No
innodb_undo_log_encrypt	Yes	Yes	Yes		Global	Yes
innodb_undo_log_truncate	Yes	Yes	Yes		Global	Yes
innodb_undo_logs	Yes	Yes	Yes		Global	Yes
innodb_undo_tablespaces	Yes	Yes	Yes		Global	Varies
innodb_use_native_aio	Yes	Yes	Yes		Global	No
innodb_version			Yes		Global	No
innodb_write_io_threads	Yes	Yes	Yes		Global	No
mecab_rc_file	Yes	Yes	Yes		Global	No
ngram_token_size	Yes	Yes	Yes		Global	No
unique_checks			Yes		Both	Yes

InnoDB Command Options

- `--ignore-builtin-innodb`

Deprecated	5.5.22	
Command-Line Format	`--ignore-builtin-innodb`	
System Variable	Name	`ignore_builtin_innodb`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean

In MySQL 5.1, this option caused the server to behave as if the built-in `InnoDB` were not present, which enabled the `InnoDB Plugin` to be used instead. In MySQL 8.0, `InnoDB` is the default storage engine and `InnoDB Plugin` is not used, so this option is ignored.

- `--innodb[=value]`

Deprecated	5.7.5	
Command-Line Format	`--innodb[=value]`	
Permitted Values	**Type**	enumeration
	Default	`ON`
	Valid Values	`OFF`
		`ON`
		`FORCE`

Controls loading of the `InnoDB` storage engine, if the server was compiled with `InnoDB` support. This option has a tristate format, with possible values of `OFF`, `ON`, or `FORCE`. See Section 5.6.2, "Installing and Uninstalling Plugins".

To disable `InnoDB`, use `--innodb=OFF` or `--skip-innodb`. In this case, because the default storage engine is `InnoDB`, the server does not start unless you also use `--default-storage-engine` and `--default-tmp-storage-engine` to set the default to some other engine for both permanent and `TEMPORARY` tables.

The `InnoDB` storage engine can no longer be disabled, and the `--innodb=OFF` and `--skip-innodb` options are deprecated and have no effect. Their use results in a warning. These options will be removed in a future MySQL release.

- `--innodb-status-file`

Command-Line Format	`--innodb-status-file`	
Permitted Values	**Type**	boolean
	Default	`OFF`

Controls whether `InnoDB` creates a file named `innodb_status.pid` in the MySQL data directory. If enabled, `InnoDB` periodically writes the output of `SHOW ENGINE INNODB STATUS` to this file.

By default, the file is not created. To create it, start `mysqld` with the `--innodb-status-file=1` option. The file is deleted during normal shutdown.

- `--skip-innodb`

Disable the `InnoDB` storage engine. See the description of `--innodb`.

InnoDB System Variables

- `daemon_memcached_enable_binlog`

Command-Line Format	`--daemon-memcached-enable-binlog=#`	
System Variable	**Name**	`daemon_memcached_enable_binlog`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`false`

Enable this option on the master server to use the `InnoDB` `memcached` plugin (`daemon_memcached`) with the MySQL binary log. This option can only be set at server startup. You must also enable the MySQL binary log on the master server using the `--log-bin` option.

For more information, see Section 15.19.7, “The InnoDB memcached Plugin and Replication”.

- `daemon_memcached_engine_lib_name`

Command-Line Format	`--daemon-memcached-engine-lib-name=library`	
System Variable	Name	`daemon_memcached_engine_lib_name`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name
	Default	`innodb_engine.so`

Specifies the shared library that implements the `InnoDB memcached` plugin.

For more information, see Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

- `daemon_memcached_engine_lib_path`

Command-Line Format	`--daemon-memcached-engine-lib-path=directory`	
System Variable	Name	`daemon_memcached_engine_lib_path`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name
	Default	`NULL`

The path of the directory containing the shared library that implements the `InnoDB memcached` plugin. The default value is NULL, representing the MySQL plugin directory. You should not need to modify this parameter unless specifying a `memcached` plugin for a different storage engine that is located outside of the MySQL plugin directory.

For more information, see Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

- `daemon_memcached_option`

Command-Line Format	`--daemon-memcached-option=options`	
System Variable	Name	`daemon_memcached_option`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	

Used to pass space-separated memcached options to the underlying `memcached` memory object caching daemon on startup. For example, you might change the port that `memcached` listens on, reduce the maximum number of simultaneous connections, change the maximum memory size for a key/value pair, or enable debugging messages for the error log.

See Section 15.19.3, “Setting Up the InnoDB memcached Plugin” for usage details. For information about `memcached` options, refer to the `memcached` man page.

- `daemon_memcached_r_batch_size`

Command-Line Format	`--daemon-memcached-r-batch-size=#`	
System Variable	Name	`daemon_memcached_r_batch_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`1`

Specifies how many `memcached` read operations (`get` operations) to perform before doing a `COMMIT` to start a new transaction. Counterpart of `daemon_memcached_w_batch_size`.

This value is set to 1 by default, so that any changes made to the table through SQL statements are immediately visible to `memcached` operations. You might increase it to reduce the overhead from frequent commits on a system where the underlying table is only being accessed through the `memcached` interface. If you set the value too large, the amount of undo or redo data could impose some storage overhead, as with any long-running transaction.

For more information, see Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

- `daemon_memcached_w_batch_size`

Command-Line Format	`--daemon-memcached-w-batch-size=#`	
System Variable	Name	`daemon_memcached_w_batch_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`1`

Specifies how many `memcached` write operations, such as `add`, `set`, and `incr`, to perform before doing a `COMMIT` to start a new transaction. Counterpart of `daemon_memcached_r_batch_size`.

This value is set to 1 by default, on the assumption that data being stored is important to preserve in case of an outage and should immediately be committed. When storing non-critical data, you might increase this value to reduce the overhead from frequent commits; but then the last *`N`*-1 uncommitted write operations could be lost if a crash occurs.

For more information, see Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

- `ignore_builtin_innodb`

Deprecated	5.5.22	
Command-Line Format	`--ignore-builtin-innodb`	
System Variable	Name	`ignore_builtin_innodb`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean

See the description of `--ignore-builtin-innodb` under “InnoDB Command Options” earlier in this section.

- `innodb_adaptive_flushing`

Command-Line Format	`--innodb-adaptive-flushing=#`	
System Variable	Name	`innodb_adaptive_flushing`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

Specifies whether to dynamically adjust the rate of flushing dirty pages in the `InnoDB` buffer pool based on the workload. Adjusting the flush rate dynamically is intended to avoid bursts of I/O activity. This setting is enabled by default. See Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing” for more information. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_adaptive_flushing_lwm`

Command-Line Format	`--innodb-adaptive-flushing-lwm=#`	
System Variable	Name	`innodb_adaptive_flushing_lwm`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`10`
	Min Value	`0`
	Max Value	`70`

Defines the low water mark representing percentage of redo log capacity at which adaptive flushing is enabled. For more information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- `innodb_adaptive_hash_index`

Command-Line Format	`--innodb-adaptive-hash-index=#`	
System Variable	Name	`innodb_adaptive_hash_index`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

Whether the `InnoDB` adaptive hash index is enabled or disabled. It may be desirable, depending on your workload, to dynamically enable or disable adaptive hash indexing to improve query

performance. Because the adaptive hash index may not be useful for all workloads, conduct benchmarks with it both enabled and disabled, using realistic workloads. See Section 15.4.3, “Adaptive Hash Index” for details.

This variable is enabled by default. You can modify this parameter using the `SET GLOBAL` statement, without restarting the server. Changing the setting requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege. You can also use `--skip-innodb_adaptive_hash_index` at server startup to disable it.

Disabling the adaptive hash index empties the hash table immediately. Normal operations can continue while the hash table is emptied, and executing queries that were using the hash table access the index B-trees directly instead. When the adaptive hash index is re-enabled, the hash table is populated again during normal operation.

- `innodb_adaptive_hash_index_parts`

Command-Line Format	`--innodb-adaptive-hash-index-parts=#`	
System Variable	Name	`innodb_adaptive_hash_index_parts`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	numeric
	Default	`8`
	Min Value	`1`
	Max Value	`512`

Partitions the adaptive hash index search system. Each index is bound to a specific partition, with each partition protected by a separate latch.

The adaptive hash index search system is partitioned into 8 parts by default. The maximum setting is 512.

For related information, see Section 15.4.3, “Adaptive Hash Index”.

- `innodb_adaptive_max_sleep_delay`

Command-Line Format	`--innodb-adaptive-max-sleep-delay=#`	
System Variable	Name	`innodb_adaptive_max_sleep_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`150000`
	Min Value	`0`
	Max Value	`1000000`

Permits `InnoDB` to automatically adjust the value of `innodb_thread_sleep_delay` up or down according to the current workload. Any non-zero value enables automated, dynamic

adjustment of the innodb_thread_sleep_delay value, up to the maximum value specified in the innodb_adaptive_max_sleep_delay option. The value represents the number of microseconds. This option can be useful in busy systems, with greater than 16 InnoDB threads. (In practice, it is most valuable for MySQL systems with hundreds or thousands of simultaneous connections.)

For more information, see Section 15.6.5, “Configuring Thread Concurrency for InnoDB”.

- innodb_api_bk_commit_interval

Command-Line Format	--innodb-api-bk-commit-interval=#	
System Variable	Name	innodb_api_bk_commit_interval
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	5
	Min Value	1
	Max Value	1073741824

How often to auto-commit idle connections that use the InnoDB memcached interface, in seconds. For more information, see Section 15.19.6.4, “Controlling Transactional Behavior of the InnoDB memcached Plugin”.

- innodb_api_disable_rowlock

Command-Line Format	--innodb-api-disable-rowlock=#	
System Variable	Name	innodb_api_disable_rowlock
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	OFF

Use this option to disable row locks when InnoDB memcached performs DML operations. By default, innodb_api_disable_rowlock is disabled, which means that memcached requests row locks for get and set operations. When innodb_api_disable_rowlock is enabled, memcached requests a table lock instead of row locks.

innodb_api_disable_rowlock is not dynamic. It must be specified on the mysqld command line or entered in the MySQL configuration file. Configuration takes effect when the plugin is installed, which occurs when the MySQL server is started.

For more information, see Section 15.19.6.4, “Controlling Transactional Behavior of the InnoDB memcached Plugin”.

- innodb_api_enable_binlog

Command-Line Format	--innodb-api-enable-binlog=#	
System Variable	Name	innodb_api_enable_binlog

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Lets you use the `InnoDB memcached` plugin with the MySQL binary log. For more information, see Enabling the InnoDB memcached Binary Log.

- `innodb_api_enable_mdl`

Command-Line Format	`--innodb-api-enable-mdl=#`	
System Variable	**Name**	`innodb_api_enable_mdl`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Locks the table used by the `InnoDB memcached` plugin, so that it cannot be dropped or altered by DDL through the SQL interface. For more information, see Section 15.19.6.4, “Controlling Transactional Behavior of the InnoDB memcached Plugin”.

- `innodb_api_trx_level`

Command-Line Format	`--innodb-api-trx-level=#`	
System Variable	**Name**	`innodb_api_trx_level`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`

Controls the transaction isolation level on queries processed by the `memcached` interface. The constants corresponding to the familiar names are:

- 0 = `READ UNCOMMITTED`
- 1 = `READ COMMITTED`
- 2 = `REPEATABLE READ`
- 3 = `SERIALIZABLE`

For more information, see Section 15.19.6.4, “Controlling Transactional Behavior of the InnoDB memcached Plugin”.

- `innodb_autoextend_increment`

Command-Line Format	`--innodb-autoextend-increment=#`	
System Variable	**Name**	`innodb_autoextend_increment`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`64`
	Min Value	`1`
	Max Value	`1000`

The increment size (in megabytes) for extending the size of an auto-extending `InnoDB` system tablespace file when it becomes full. The default value is 64. For related information, see System Tablespace Data File Configuration, and Section 15.7.1, “Resizing the InnoDB System Tablespace”.

The `innodb_autoextend_increment` setting does not affect file-per-table tablespace files or general tablespace files. These files are auto-extending regardless of the `innodb_autoextend_increment` setting. The initial extensions are by small amounts, after which extensions occur in increments of 4MB.

- `innodb_autoinc_lock_mode`

Command-Line Format	`--innodb-autoinc-lock-mode=#`	
System Variable	**Name**	`innodb_autoinc_lock_mode`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`1`
	Valid Values	`0`
		`1`
		`2`

The lock mode to use for generating auto-increment values. Permissible values are 0, 1, or 2, for “traditional”, “consecutive”, or “interleaved”, respectively. The default setting is 1 (“consecutive”). For the characteristics of each lock mode, see InnoDB AUTO_INCREMENT Lock Modes.

- `innodb_background_drop_list_empty`

Command-Line Format	`--innodb-background-drop-list-empty=#`	
System Variable	**Name**	`innodb_background_drop_list_empty`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Enabling the `innodb_background_drop_list_empty` debug option helps avoid test case failures by delaying table creation until the background drop list is empty. For example, if test case A

places table `t1` on the background drop list, test case B waits until the background drop list is empty before creating table `t1`.

- `innodb_buffer_pool_chunk_size`

Command-Line Format	`--innodb-buffer-pool-chunk-size`	
System Variable	Name	`innodb_buffer_pool_chunk_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`134217728`
	Min Value	`1048576`
	Max Value	`innodb_buffer_pool_size / innodb_buffer_pool_instances`

`innodb_buffer_pool_chunk_size` defines the chunk size for `InnoDB` buffer pool resizing operations. The `innodb_buffer_pool_size` parameter is dynamic, which allows you to resize the buffer pool without restarting the server.

To avoid copying all buffer pool pages during resizing operations, the operation is performed in "chunks". By default, `innodb_buffer_pool_chunk_size` is 128MB (134217728 bytes). The number of pages contained in a chunk depends on the value of `innodb_page_size`. `innodb_buffer_pool_chunk_size` can be increased or decreased in units of 1MB (1048576 bytes).

The following conditions apply when altering the `innodb_buffer_pool_chunk_size` value:

- If `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` is larger than the current buffer pool size when the buffer pool is initialized, `innodb_buffer_pool_chunk_size` is truncated to `innodb_buffer_pool_size` / `innodb_buffer_pool_instances`.
- Buffer pool size must always be equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances`. If you alter `innodb_buffer_pool_chunk_size`, `innodb_buffer_pool_size` is automatically adjusted to a value that is equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` that is not less than the current buffer pool size. The adjustment occurs when the buffer pool is initialized.

Important

Care should be taken when changing `innodb_buffer_pool_chunk_size`, as changing this value can automatically increase the size of the buffer pool. Before changing `innodb_buffer_pool_chunk_size`, calculate the effect it will have on `innodb_buffer_pool_size` to ensure that the resulting buffer pool size is acceptable.

To avoid potential performance issues, the number of chunks (`innodb_buffer_pool_size` / `innodb_buffer_pool_chunk_size`) should not exceed 1000.

See Section 15.6.3.2, "Configuring InnoDB Buffer Pool Size" for more information.

- `innodb_buffer_pool_debug`

Command-Line Format	`--innodb-buffer-pool-debug=#`

System Variable	Name	innodb_buffer_pool_debug
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	OFF

Enabling this option permits multiple buffer pool instances when the buffer pool is less than 1GB in size, ignoring the 1GB minimum buffer pool size constraint imposed on innodb_buffer_pool_instances. The innodb_buffer_pool_debug option is only available if debugging support is compiled in using the WITH_DEBUG CMake option.

- innodb_buffer_pool_dump_at_shutdown

Command-Line Format	--innodb-buffer-pool-dump-at-shutdown=#	
System Variable	Name	innodb_buffer_pool_dump_at_shutdown
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	ON

Specifies whether to record the pages cached in the InnoDB buffer pool when the MySQL server is shut down, to shorten the warmup process at the next restart. Typically used in combination with innodb_buffer_pool_load_at_startup. The innodb_buffer_pool_dump_pct option defines the percentage of most recently used buffer pool pages to dump.

Both innodb_buffer_pool_dump_at_shutdown and innodb_buffer_pool_load_at_startup are enabled by default.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- innodb_buffer_pool_dump_now

Command-Line Format	--innodb-buffer-pool-dump-now=#	
System Variable	Name	innodb_buffer_pool_dump_now
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Immediately records the pages cached in the InnoDB buffer pool. Typically used in combination with innodb_buffer_pool_load_now.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- innodb_buffer_pool_dump_pct

Command-Line Format	--innodb-buffer-pool-dump-pct=#

System Variable	Name	innodb_buffer_pool_dump_pct
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	25
	Min Value	1
	Max Value	100

Specifies the percentage of the most recently used pages for each buffer pool to read out and dump. The range is 1 to 100. The default value is 25. For example, if there are 4 buffer pools with 100 pages each, and `innodb_buffer_pool_dump_pct` is set to 25, the 25 most recently used pages from each buffer pool are dumped.

- `innodb_buffer_pool_filename`

Command-Line Format	--innodb-buffer-pool-filename=file	
System Variable	Name	innodb_buffer_pool_filename
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	file name
	Default	ib_buffer_pool

Specifies the name of the file that holds the list of tablespace IDs and page IDs produced by `innodb_buffer_pool_dump_at_shutdown` or `innodb_buffer_pool_dump_now`. Tablespace IDs and page IDs are saved in the following format: `space, page_id`. By default, the file is named `ib_buffer_pool` and is located in the `InnoDB` data directory. A non-default location must be specified relative to the data directory.

A file name can be specified at runtime, using a `SET` statement:

```
SET GLOBAL innodb_buffer_pool_filename='file_name';
```

You can also specify a file name at startup, in a startup string or MySQL configuration file. When specifying a file name at startup, the file must exist or `InnoDB` will return a startup error indicating that there is no such file or directory.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `innodb_buffer_pool_instances`

Command-Line Format	--innodb-buffer-pool-instances=#	
System Variable	Name	innodb_buffer_pool_instances
	Variable Scope	Global
	Dynamic Variable	No

Permitted Values (Windows, 32-bit platforms)	**Type**	integer
	Default	`(autosized)`
	Min Value	`1`
	Max Value	`64`
Permitted Values (Other)	**Type**	integer
	Default	`8 (or 1 if innodb_buffer_pool_size < 1GB`
	Min Value	`1`
	Max Value	`64`

The number of regions that the `InnoDB` buffer pool is divided into. For systems with buffer pools in the multi-gigabyte range, dividing the buffer pool into separate instances can improve concurrency, by reducing contention as different threads read and write to cached pages. Each page that is stored in or read from the buffer pool is assigned to one of the buffer pool instances randomly, using a hashing function. Each buffer pool manages its own free lists, flush lists, LRUs, and all other data structures connected to a buffer pool, and is protected by its own buffer pool mutex.

This option only takes effect when setting `innodb_buffer_pool_size` to 1GB or more. The total buffer pool size is divided among all the buffer pools. For best efficiency, specify a combination of `innodb_buffer_pool_instances` and `innodb_buffer_pool_size` so that each buffer pool instance is at least 1GB.

The default value on 32-bit Windows systems depends on the value of `innodb_buffer_pool_size`, as described below:

- If `innodb_buffer_pool_size` is greater than 1.3GB, the default for `innodb_buffer_pool_instances` is `innodb_buffer_pool_size`/128MB, with individual memory allocation requests for each chunk. 1.3GB was chosen as the boundary at which there is significant risk for 32-bit Windows to be unable to allocate the contiguous address space needed for a single buffer pool.
- Otherwise, the default is 1.

On all other platforms, the default value is 8 when `innodb_buffer_pool_size` is greater than or equal to 1GB. Otherwise, the default is 1.

For related information, see Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size”.

- `innodb_buffer_pool_load_abort`

Command-Line Format	`--innodb-buffer-pool-load-abort=#`	
System Variable	**Name**	`innodb_buffer_pool_load_abort`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Interrupts the process of restoring `InnoDB` buffer pool contents triggered by `innodb_buffer_pool_load_at_startup` or `innodb_buffer_pool_load_now`.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `innodb_buffer_pool_load_at_startup`

Command-Line Format	`--innodb-buffer-pool-load-at-startup=#`	
System Variable	Name	`innodb_buffer_pool_load_at_startup`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`ON`

Specifies that, on MySQL server startup, the `InnoDB` buffer pool is automatically warmed up by loading the same pages it held at an earlier time. Typically used in combination with `innodb_buffer_pool_dump_at_shutdown`.

Both `innodb_buffer_pool_dump_at_shutdown` and `innodb_buffer_pool_load_at_startup` are enabled by default.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `innodb_buffer_pool_load_now`

Command-Line Format	`--innodb-buffer-pool-load-now=#`	
System Variable	Name	`innodb_buffer_pool_load_now`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Immediately warms up the `InnoDB` buffer pool by loading a set of data pages, without waiting for a server restart. Can be useful to bring cache memory back to a known state during benchmarking, or to ready the MySQL server to resume its normal workload after running queries for reports or maintenance.

For more information, see Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

- `innodb_buffer_pool_size`

Command-Line Format	`--innodb-buffer-pool-size=#`	
System Variable	Name	`innodb_buffer_pool_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	`134217728`
	Min Value	`5242880`

	Max Value	`2**32-1`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`134217728`
	Min Value	`5242880`
	Max Value	`2**64-1`

The size in bytes of the buffer pool, the memory area where `InnoDB` caches table and index data. The default value is 128MB. The maximum value depends on the CPU architecture; the maximum is 4294967295 (2^{32}-1) on 32-bit systems and 18446744073709551615 (2^{64}-1) on 64-bit systems. On 32-bit systems, the CPU architecture and operating system may impose a lower practical maximum size than the stated maximum. When the size of the buffer pool is greater than 1GB, setting `innodb_buffer_pool_instances` to a value greater than 1 can improve the scalability on a busy server.

A larger buffer pool requires less disk I/O to access the same table data more than once. On a dedicated database server, you might set the buffer pool size to 80% of the machine's physical memory size. Be aware of the following potential issues when configuring buffer pool size, and be prepared to scale back the size of the buffer pool if necessary.

- Competition for physical memory can cause paging in the operating system.
- `InnoDB` reserves additional memory for buffers and control structures, so that the total allocated space is approximately 10% greater than the specified buffer pool size.
- Address space for the buffer pool must be contiguous, which can be an issue on Windows systems with DLLs that load at specific addresses.
- The time to initialize the buffer pool is roughly proportional to its size. On instances with large buffer pools, initialization time might be significant. To reduce the initialization period, you can save the buffer pool state at server shutdown and restore it at server startup. See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.

When you increase or decrease buffer pool size, the operation is performed in chunks. Chunk size is defined by the `innodb_buffer_pool_chunk_size` configuration option, which has a default of 128 MB.

Buffer pool size must always be equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances`. If you alter the buffer pool size to a value that is not equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances`, buffer pool size is automatically adjusted to a value that is equal to or a multiple of `innodb_buffer_pool_chunk_size` * `innodb_buffer_pool_instances` that is not less than the specified buffer pool size.

`innodb_buffer_pool_size` can be set dynamically, which allows you to resize the buffer pool without restarting the server. The `Innodb_buffer_pool_resize_status` status variable reports the status of online buffer pool resizing operations. See Section 15.6.3.2, “Configuring InnoDB Buffer Pool Size” for more information.

- `innodb_change_buffer_max_size`

Command-Line Format	`--innodb-change-buffer-max-size=#`	
System Variable	**Name**	`innodb_change_buffer_max_size`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`25`
	Min Value	`0`
	Max Value	`50`

Maximum size for the `InnoDB` change buffer, as a percentage of the total size of the buffer pool. You might increase this value for a MySQL server with heavy insert, update, and delete activity, or decrease it for a MySQL server with unchanging data used for reporting. For more information, see Section 15.4.2, “Change Buffer”, and Section 15.6.4, “Configuring InnoDB Change Buffering”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_change_buffering`

Command-Line Format	`--innodb-change-buffering=#`	
System Variable	**Name**	`innodb_change_buffering`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`all`
	Valid Values	`none`
		`inserts`
		`deletes`
		`changes`
		`purges`
		`all`

Whether `InnoDB` performs change buffering, an optimization that delays write operations to secondary indexes so that the I/O operations can be performed sequentially. Permitted values are described in the following table. Values may also be specified numerically.

Table 15.13 Permitted Values for innodb_change_buffering

Value	Numeric Value	Description
`none`	`0`	Do not buffer any operations.
`inserts`	`1`	Buffer insert operations.
`deletes`	`2`	Buffer delete marking operations; strictly speaking, the writes that mark index records for later deletion during a purge operation.
`changes`	`3`	Buffer inserts and delete-marking operations.
`purges`	`4`	Buffer the physical deletion operations that happen in the background.
`all`	`5`	The default. Buffer inserts, delete-marking operations, and purges.

For more information, see Section 15.4.2, “Change Buffer”, and Section 15.6.4, “Configuring InnoDB Change Buffering”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_change_buffering_debug`

Command-Line Format	--innodb-change-buffering-debug=#	
System Variable	Name	innodb_change_buffering_debug
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	0
	Max Value	2

Sets a debug flag for InnoDB change buffering. A value of 1 forces all changes to the change buffer. A value of 2 causes a crash at merge. A default value of 0 indicates that the change buffering debug flag is not set. This option is only available when debugging support is compiled in using the WITH_DEBUG CMake option.

- `innodb_checkpoint_disabled`

Introduced	8.0.2	
Command-Line Format	--innodb-checkpoint-disabled=#	
System Variable	Name	innodb_checkpoint_disabled
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

This is a debug option that is only intended for expert debugging use. It disables checkpoints so that a deliberate server exit always initiates InnoDB recovery. It should only be enabled for a short interval, typically before running DML operations that write redo log entries that would require recovery following a server exit. This option is only available if debugging support is compiled in using the WITH_DEBUG CMake option.

- `innodb_checksum_algorithm`

Command-Line Format	--innodb-checksum-algorithm=#	
System Variable	Name	innodb_checksum_algorithm
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	crc32
	Valid Values	innodb

		crc32
		none
		strict_innodb
		strict_crc32
		strict_none

Specifies how to generate and verify the checksum stored in the disk blocks of InnoDB tablespaces. The default value for innodb_checksum_algorithm is crc32.

Versions of MySQL Enterprise Backup up to 3.8.0 do not support backing up tablespaces that use CRC32 checksums. MySQL Enterprise Backup adds CRC32 checksum support in 3.8.1, with some limitations. Refer to the MySQL Enterprise Backup 3.8.1 Change History for more information.

The value innodb is backward-compatible with earlier versions of MySQL. The value crc32 uses an algorithm that is faster to compute the checksum for every modified block, and to check the checksums for each disk read. It scans blocks 32 bits at a time, which is faster than the innodb checksum algorithm, which scans blocks 8 bits at a time. The value none writes a constant value in the checksum field rather than computing a value based on the block data. The blocks in a tablespace can use a mix of old, new, and no checksum values, being updated gradually as the data is modified; once blocks in a tablespace are modified to use the crc32 algorithm, the associated tables cannot be read by earlier versions of MySQL.

The strict form of a checksum algorithm reports an error if it encounters a valid but non-matching checksum value in a tablespace. It is recommended that you only use strict settings in a new instance, to set up tablespaces for the first time. Strict settings are somewhat faster, because they do not need to compute all checksum values during disk reads.

The following table shows the difference between the none, innodb, and crc32 option values, and their strict counterparts. none, innodb, and crc32 write the specified type of checksum value into each data block, but for compatibility accept other checksum values when verifying a block during a read operation. Strict settings also accept valid checksum values but print an error message when a valid non-matching checksum value is encountered. Using the strict form can make verification faster if all InnoDB data files in an instance are created under an identical innodb_checksum_algorithm value.

Table 15.14 innodb_checksum_algorithm Settings

Value	Generated checksum (when writing)	Permitted checksums (when reading)
none	A constant number.	Any of the checksums generated by none, innodb, or crc32.
innodb	A checksum calculated in software, using the original algorithm from InnoDB.	Any of the checksums generated by none, innodb, or crc32.
crc32	A checksum calculated using the crc32 algorithm, possibly done with a hardware assist.	Any of the checksums generated by none, innodb, or crc32.
strict_none	A constant number	Any of the checksums generated by none, innodb, or crc32. InnoDB prints an error message if a valid but non-matching checksum is encountered.
strict_innodb	A checksum calculated in software, using the original algorithm from InnoDB.	Any of the checksums generated by none, innodb, or crc32. InnoDB prints an error message if a valid

Value	Generated checksum (when writing)	Permitted checksums (when reading)
		but non-matching checksum is encountered.
strict_crc32	A checksum calculated using the `crc32` algorithm, possibly done with a hardware assist.	Any of the checksums generated by `none`, `innodb`, or `crc32`. `InnoDB` prints an error message if a valid but non-matching checksum is encountered.

- `innodb_cmp_per_index_enabled`

Command-Line Format	`--innodb-cmp-per-index-enabled=#`	
System Variable	Name	`innodb_cmp_per_index_enabled`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`
	Valid Values	`OFF`
		`ON`

Enables per-index compression-related statistics in the `INFORMATION_SCHEMA.INNODB_CMP_PER_INDEX` table. Because these statistics can be expensive to gather, only enable this option on development, test, or slave instances during performance tuning related to `InnoDB` compressed tables.

For more information, see Section 24.31.7, “The INFORMATION_SCHEMA INNODB_CMP_PER_INDEX and INNODB_CMP_PER_INDEX_RESET Tables”, and Section 15.9.1.4, “Monitoring InnoDB Table Compression at Runtime”.

- `innodb_commit_concurrency`

Command-Line Format	`--innodb-commit-concurrency=#`	
System Variable	Name	`innodb_commit_concurrency`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1000`

The number of threads that can commit at the same time. A value of 0 (the default) permits any number of transactions to commit simultaneously.

The value of `innodb_commit_concurrency` cannot be changed at runtime from zero to nonzero or vice versa. The value can be changed from one nonzero value to another.

- `innodb_compress_debug`

Command-Line Format	`--innodb-compress-debug=#`	
System Variable	Name	`innodb_compress_debug`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`none`
	Valid Values	`none`
		`zlib`
		`lz4`
		`lz4hc`

Compresses all tables using a specified compression algorithm without having to define a `COMPRESSION` attribute for each table. This option is only available if debugging support is compiled in using the `WITH_DEBUG CMake` option.

For related information, see Section 15.9.2, “InnoDB Page Compression”.

- `innodb_compression_failure_threshold_pct`

Command-Line Format	`--innodb-compression-failure-threshold-pct=#`	
System Variable	Name	`innodb_compression_failure_threshold_pct`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`5`
	Min Value	`0`
	Max Value	`100`

Sets the cutoff point at which MySQL begins adding padding within compressed pages to avoid expensive compression failures. A value of zero disables the mechanism that monitors compression efficiency and dynamically adjusts the padding amount.

For more information, see Section 15.9.1.6, “Compression for OLTP Workloads”.

- `innodb_compression_level`

Command-Line Format	`--innodb-compression-level=#`	
System Variable	Name	`innodb_compression_level`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer

	Default	6
	Min Value	0
	Max Value	9

Specifies the level of zlib compression to use for InnoDB compressed tables and indexes.

For more information, see Section 15.9.1.6, “Compression for OLTP Workloads”.

- innodb_compression_pad_pct_max

Command-Line Format	--innodb-compression-pad-pct-max=#	
System Variable	**Name**	innodb_compression_pad_pct_max
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	50
	Min Value	0
	Max Value	75

Specifies the maximum percentage that can be reserved as free space within each compressed page, allowing room to reorganize the data and modification log within the page when a compressed table or index is updated and the data might be recompressed. Only applies when innodb_compression_failure_threshold_pct is set to a non-zero value, and the rate of compression failures passes the cutoff point.

For more information, see Section 15.9.1.6, “Compression for OLTP Workloads”.

- innodb_concurrency_tickets

Command-Line Format	--innodb-concurrency-tickets=#	
System Variable	**Name**	innodb_concurrency_tickets
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	5000
	Min Value	1
	Max Value	4294967295

Determines the number of threads that can enter InnoDB concurrently. A thread is placed in a queue when it tries to enter InnoDB if the number of threads has already reached the concurrency limit. When a thread is permitted to enter InnoDB, it is given a number of “ tickets” equal to the value of innodb_concurrency_tickets, and the thread can enter and leave InnoDB freely until it has

used up its tickets. After that point, the thread again becomes subject to the concurrency check (and possible queuing) the next time it tries to enter InnoDB. The default value is 5000.

With a small innodb_concurrency_tickets value, small transactions that only need to process a few rows compete fairly with larger transactions that process many rows. The disadvantage of a small innodb_concurrency_tickets value is that large transactions must loop through the queue many times before they can complete, which extends the amount of time required to complete their task.

With a large innodb_concurrency_tickets value, large transactions spend less time waiting for a position at the end of the queue (controlled by innodb_thread_concurrency) and more time retrieving rows. Large transactions also require fewer trips through the queue to complete their task. The disadvantage of a large innodb_concurrency_tickets value is that too many large transactions running at the same time can starve smaller transactions by making them wait a longer time before executing.

With a non-zero innodb_thread_concurrency value, you may need to adjust the innodb_concurrency_tickets value up or down to find the optimal balance between larger and smaller transactions. The SHOW ENGINE INNODB STATUS report shows the number of tickets remaining for an executing transaction in its current pass through the queue. This data may also be obtained from the TRX_CONCURRENCY_TICKETS column of the INFORMATION_SCHEMA.INNODB_TRX table.

For more information, see Section 15.6.5, “Configuring Thread Concurrency for InnoDB”.

- innodb_data_file_path

Command-Line Format	--innodb-data-file-path=name	
System Variable	Name	innodb_data_file_path
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string
	Default	ibdata1:12M:autoextend

Defines the path and file size for individual InnoDB system tablespace data files. The full directory path for system tablespace data files is formed by concatenating path defined by innodb_data_home_dir and innodb_data_file_path. File sizes are specified KB, MB or GB (1024MB) by appending K, M or G to the size value. If specifying the data file size in kilobytes (KB), do so in multiples of 1024. Otherwise, KB values are rounded to nearest megabyte (MB) boundary. The sum of the sizes of the files must be at least slightly larger than 10MB. If you do not specify innodb_data_file_path, the default behavior is to create a single auto-extending data file, slightly larger than 12MB, named ibdata1. The size limit of individual files is determined by your operating system. You can set the file size to more than 4GB on operating systems that support large files. You can also use raw disk partitions as data files. For more information about configuring system tablespace data files, see Section 15.6.1, “InnoDB Startup Configuration”.

The following minimum file sizes are enforced for the *first* system tablespace data file (ibdata1) to ensure that there is enough space for doublewrite buffer pages:

- For an innodb_page_size value of 16KB or less, the minimum data file size is 3MB.
- For an innodb_page_size value of 32KB, the minimum data file size is 6MB.
- For an innodb_page_size value of 64KB, the minimum data file size is 12MB.

- innodb_data_home_dir

Command-Line Format	--innodb-data-home-dir=dir_name	
System Variable	Name	innodb_data_home_dir
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The common part of the directory path for InnoDB system tablespace data files. This setting does not affect the location of file-per-table tablespaces when innodb_file_per_table is enabled. The default value is the MySQL data directory. If you specify the value as an empty string, you can specify an absolute file paths for innodb_data_file_path.

A trailing slash is required when specifying a value for innodb_data_home_dir. For example:

```
[mysqld]
innodb_data_home_dir = /path/to/myibdata/
```

For related information, see Section 15.6.1, “InnoDB Startup Configuration”.

- innodb_deadlock_detect

Command-Line Format	--innodb-deadlock-detect	
System Variable	Name	innodb_deadlock_detect
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	ON

This option is used to disable deadlock detection. On high concurrency systems, deadlock detection can cause a slowdown when numerous threads wait for the same lock. At times, it may be more efficient to disable deadlock detection and rely on the innodb_lock_wait_timeout setting for transaction rollback when a deadlock occurs.

For related information, see Section 15.5.5.2, “Deadlock Detection and Rollback”.

- innodb_default_row_format

Command-Line Format	--innodb-default-row-format=#	
System Variable	Name	innodb_default_row_format
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	DYNAMIC
	Valid Values	DYNAMIC
		COMPACT
		REDUNDANT

The innodb_default_row_format option defines the default row format for InnoDB tables and user-created temporary tables. The default setting is DYNAMIC. Other permitted values are COMPACT and REDUNDANT. The COMPRESSED row format, which is not supported for use in the system tablespace, cannot be defined as the default.

Newly created tables use the row format defined by innodb_default_row_format when a ROW_FORMAT option is not specified explicitly or when ROW_FORMAT=DEFAULT is used.

When a ROW_FORMAT option is not specified explicitly or when ROW_FORMAT=DEFAULT is used, any operation that rebuilds a table also silently changes the row format of the table to the format defined by innodb_default_row_format. For more information, see Section 15.10.2, "Specifying the Row Format for a Table".

Internal InnoDB temporary tables created by the server to process queries use the DYNAMIC row format, regardless of the innodb_default_row_format setting.

- innodb_disable_sort_file_cache

Command-Line Format	--innodb-disable-sort-file-cache=#	
System Variable	Name	innodb_disable_sort_file_cache
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Disables the operating system file system cache for merge-sort temporary files. The effect is to open such files with the equivalent of O_DIRECT.

- innodb_doublewrite

Command-Line Format	--innodb-doublewrite	
System Variable	Name	innodb_doublewrite
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	ON

When enabled (the default), InnoDB stores all data twice, first to the doublewrite buffer, then to the actual data files. This variable can be turned off with --skip-innodb_doublewrite for benchmarks or cases when top performance is needed rather than concern for data integrity or possible failures.

If system tablespace data files (ibdata* files) are located on Fusion-io devices that support atomic writes, doublewrite buffering is automatically disabled and Fusion-io atomic writes are used for all data files. Because the doublewrite buffer setting is global, doublewrite buffering is also disabled for data files residing on non-Fusion-io hardware. This feature is only supported on Fusion-io hardware and only enabled for Fusion-io NVMFS on Linux. To take full advantage of this feature, an innodb_flush_method setting of O_DIRECT is recommended.

For related information, see Section 15.4.6, "Doublewrite Buffer".

- `innodb_fast_shutdown`

Command-Line Format	`--innodb-fast-shutdown[=#]`	
System Variable	Name	`innodb_fast_shutdown`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1`
	Valid Values	`0`
		`1`
		`2`

The `InnoDB` shutdown mode. If the value is 0, `InnoDB` does a slow shutdown, a full purge and a change buffer merge before shutting down. If the value is 1 (the default), `InnoDB` skips these operations at shutdown, a process known as a fast shutdown. If the value is 2, `InnoDB` flushes its logs and shuts down cold, as if MySQL had crashed; no committed transactions are lost, but the crash recovery operation makes the next startup take longer.

The slow shutdown can take minutes, or even hours in extreme cases where substantial amounts of data are still buffered. Use the slow shutdown technique before upgrading or downgrading between MySQL major releases, so that all data files are fully prepared in case the upgrade process updates the file format.

Use `innodb_fast_shutdown=2` in emergency or troubleshooting situations, to get the absolute fastest shutdown if data is at risk of corruption.

- `innodb_fil_make_page_dirty_debug`

Command-Line Format	`--innodb-fil-make-page-dirty-debug=#`	
System Variable	Name	`innodb_fil_make_page_dirty_debug`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Max Value	`2**32-1`

By default, setting `innodb_fil_make_page_dirty_debug` to the ID of a tablespace immediately dirties the first page of the tablespace. If `innodb_saved_page_number_debug` is set to a non-default value, setting `innodb_fil_make_page_dirty_debug` dirties the specified page. The `innodb_fil_make_page_dirty_debug` option is only available if debugging support is compiled in using the `WITH_DEBUG` `CMake` option.

- `innodb_file_per_table`

Command-Line Format	`--innodb-file-per-table`	
System Variable	Name	`innodb_file_per_table`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

When `innodb_file_per_table` is enabled (the default), `InnoDB` stores the data and indexes for each newly created table in a separate `.ibd` file instead of the system tablespace. The storage for these tables is reclaimed when the tables are dropped or truncated. This setting enables `InnoDB`features such as table compression. See Section 15.7.4, “InnoDB File-Per-Table Tablespaces” for more information.

Enabling `innodb_file_per_table` also means that an `ALTER TABLE` operation moves an `InnoDB` table from the system tablespace to an individual `.ibd` file in cases where `ALTER TABLE` rebuilds the table (`ALGORITHM=COPY`). An exception to this rule is for tables placed in the system tablespace using the `TABLESPACE=innodb_system` option with `CREATE TABLE` or `ALTER TABLE`. These tables are unaffected by the `innodb_file_per_table` setting and can only be moved to file-per-table tablespaces using `ALTER TABLE ... TABLESPACE=innodb_file_per_table`.

When `innodb_file_per_table` is disabled, `InnoDB` stores the data for tables and indexes in the ibdata files that make up the system tablespace. This setting reduces the performance overhead of file system operations for operations such as `DROP TABLE` or `TRUNCATE TABLE`. It is most appropriate for a server environment where entire storage devices are devoted to MySQL data. Because the system tablespace never shrinks, and is shared across all databases in an instance, avoid loading huge amounts of temporary data on a space-constrained system when `innodb_file_per_table` is disabled. Set up a separate instance in such cases, so that you can drop the entire instance to reclaim the space.

`innodb_file_per_table` is enabled by default. Consider disabling it if backward compatibility with MySQL 5.5 or earlier is a concern. This will prevent `ALTER TABLE` from moving `InnoDB` tables from the system tablespace to individual `.ibd` files.

`innodb_file_per_table` is dynamic and can be set `ON` or `OFF` using `SET GLOBAL`. You can also set this option in the MySQL configuration file (`my.cnf` or `my.ini`) but this requires shutting down and restarting the server.

Dynamically changing the value requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege and immediately affects the operation of all connections.

The `innodb_file_per-table` setting does not affect the creation of `InnoDB` temporary tables. All `InnoDB` temporary tables are created in the shared temporary tablespace.

- `innodb_fill_factor`

Command-Line Format	`--innodb-fill-factor=#`	
System Variable	**Name**	`innodb_fill_factor`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`100`
	Min Value	`10`

	Max Value	100

InnoDB performs a bulk load when creating or rebuilding indexes. This method of index creation is known as a “sorted index build”.

innodb_fill_factor defines the percentage of space on each B-tree page that is filled during a sorted index build, with the remaining space reserved for future index growth. For example, setting innodb_fill_factor to 80 reserves 20 percent of the space on each B-tree page for future index growth. Actual percentages may vary. The innodb_fill_factor setting is interpreted as a hint rather than a hard limit.

An innodb_fill_factor setting of 100 leaves 1/16 of the space in clustered index pages free for future index growth.

innodb_fill_factor applies to both B-tree leaf and non-leaf pages. It does not apply to external pages used for TEXT or BLOB entries.

For more information, see Section 15.8.2.3, “Sorted Index Builds”.

- innodb_flush_log_at_timeout

System Variable	**Name**	innodb_flush_log_at_timeout
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	1
	Min Value	1
	Max Value	2700

Write and flush the logs every *N* seconds. innodb_flush_log_at_timeout allows the timeout period between flushes to be increased in order to reduce flushing and avoid impacting performance of binary log group commit. The default setting for innodb_flush_log_at_timeout is once per second.

- innodb_flush_log_at_trx_commit

Command-Line Format	--innodb-flush-log-at-trx-commit[=#]	
System Variable	**Name**	innodb_flush_log_at_trx_commit
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	1
	Valid Values	0
		1
		2

Controls the balance between strict ACID compliance for commit operations and higher performance that is possible when commit-related I/O operations are rearranged and done in batches. You can achieve better performance by changing the default value but then you can lose up to a second of transactions in a crash.

- The default value of 1 is required for full ACID compliance. With this value, the contents of the `InnoDB` log buffer are written out to the log file at each transaction commit and the log file is flushed to disk.
- With a value of 0, the contents of the `InnoDB` log buffer are written to the log file approximately once per second and the log file is flushed to disk. No writes from the log buffer to the log file are performed at transaction commit. Once-per-second flushing is not 100% guaranteed to happen every second, due to process scheduling issues. Because the flush to disk operation only occurs approximately once per second, you can lose up to a second of transactions with any `mysqld` process crash.
- With a value of 2, the contents of the `InnoDB` log buffer are written to the log file after each transaction commit and the log file is flushed to disk approximately once per second. Once-per-second flushing is not guaranteed to happen every second due to process scheduling issues. Because the flush to disk operation only occurs approximately once per second, you can lose up to a second of transactions in an operating system crash or a power outage.
- `InnoDB` log flushing frequency is controlled by `innodb_flush_log_at_timeout`, which allows you to set log flushing frequency to *`N`* seconds (where *`N`* is `1 ... 2700`, with a default value of 1). However, any `mysqld` process crash can erase up to *`N`* seconds of transactions.
- DDL changes and other internal `InnoDB` activities flush the `InnoDB` log independent of the `innodb_flush_log_at_trx_commit` setting.
- `InnoDB` crash recovery works regardless of the `innodb_flush_log_at_trx_commit` setting. Transactions are either applied entirely or erased entirely.

For durability and consistency in a replication setup that uses `InnoDB` with transactions:

- If binary logging is enabled, set `sync_binlog=1`.
- Always set `innodb_flush_log_at_trx_commit=1`.

Caution

Many operating systems and some disk hardware fool the flush-to-disk operation. They may tell `mysqld` that the flush has taken place, even though it has not. In this case, the durability of transactions is not guaranteed even with the setting 1, and in the worst case, a power outage can corrupt `InnoDB` data. Using a battery-backed disk cache in the SCSI disk controller or in the disk itself speeds up file flushes, and makes the operation safer. You can also try to disable the caching of disk writes in hardware caches.

- `innodb_flush_method`

Command-Line Format	`--innodb-flush-method=name`	
System Variable	**Name**	`innodb_flush_method`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (Unix)	**Type**	string

	Default	`fsync`
	Valid Values	`fsync`
		`O_DSYNC`
		`littlesync`
		`nosync`
		`O_DIRECT`
		`O_DIRECT_NO_FSYNC`
Permitted Values (Windows)	**Type**	string
	Default	`unbuffered`
	Valid Values	`unbuffered`
		`normal`

Defines the method used to flush data to `InnoDB` data files and log files, which can affect I/O throughput.

On Unix-like systems, the default value is `fsync`. On Windows, the default value is `unbuffered`.

Note

In MySQL 8.0, `innodb_flush_method` options may be specified numerically.

The `innodb_flush_method` options for Unix-like systems include:

- `fsync` or `0`: `InnoDB` uses the `fsync()` system call to flush both the data and log files. `fsync` is the default setting.
- `O_DSYNC` or `1`: `InnoDB` uses `O_SYNC` to open and flush the log files, and `fsync()` to flush the data files. `InnoDB` does not use `O_DSYNC` directly because there have been problems with it on many varieties of Unix.
- `littlesync` or `2`: This option is used for internal performance testing and is currently unsupported. Use at your own risk.
- `nosync` or `3`: This option is used for internal performance testing and is currently unsupported. Use at your own risk.
- `O_DIRECT` or `4`: `InnoDB` uses `O_DIRECT` (or `directio()` on Solaris) to open the data files, and uses `fsync()` to flush both the data and log files. This option is available on some GNU/Linux versions, FreeBSD, and Solaris.
- `O_DIRECT_NO_FSYNC` or `5`: `InnoDB` uses `O_DIRECT` during flushing I/O, but skips the `fsync()` system call afterward. This setting is suitable for some types of file systems but not others. For example, it is not suitable for XFS. If you are not sure whether the file system you use requires an `fsync()`, for example to preserve all file metadata, use `O_DIRECT` instead.

The `innodb_flush_method` options for Windows systems include:

- `unbuffered` or `0`: `InnoDB` uses simulated asynchronous I/O and non-buffered I/O.
- `normal` or `1`: `InnoDB` uses simulated asynchronous I/O and buffered I/O.

How each setting affects performance depends on hardware configuration and workload. Benchmark your particular configuration to decide which setting to use, or whether to keep the default setting. Examine the `Innodb_data_fsyncs` status variable to see the overall number of `fsync()` calls for each setting. The mix of read and write operations in your workload can affect how a setting

performs. For example, on a system with a hardware RAID controller and battery-backed write cache, `O_DIRECT` can help to avoid double buffering between the `InnoDB` buffer pool and the operating system file system cache. On some systems where `InnoDB` data and log files are located on a SAN, the default value or `O_DSYNC` might be faster for a read-heavy workload with mostly `SELECT` statements. Always test this parameter with hardware and workload that reflect your production environment. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_flush_neighbors`

Command-Line Format	`--innodb-flush-neighbors`	
System Variable	Name	`innodb_flush_neighbors`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`1`
	Valid Values	`0`
		`1`
		`2`

Specifies whether flushing a page from the `InnoDB` buffer pool also flushes other dirty pages in the same extent.

- The default value of 1 flushes contiguous dirty pages in the same extent from the buffer pool.
- A setting of 0 turns `innodb_flush_neighbors` off and no other dirty pages are flushed from the buffer pool.
- A setting of 2 flushes dirty pages in the same extent from the buffer pool.

When the table data is stored on a traditional HDD storage device, flushing such neighbor pages in one operation reduces I/O overhead (primarily for disk seek operations) compared to flushing individual pages at different times. For table data stored on SSD, seek time is not a significant factor and you can turn this setting off to spread out write operations. For related information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- `innodb_flush_sync`

Command-Line Format	`--innodb-flush-sync=#`	
System Variable	Name	`innodb_flush_sync`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

The `innodb_flush_sync` parameter, which is enabled by default, causes the `innodb_io_capacity` setting to be ignored for bursts of I/O activity that occur at checkpoints. To adhere to the limit on `InnoDB` background I/O activity defined by the `innodb_io_capacity` setting, disable `innodb_flush_sync`.

For related information, see Section 15.6.8, “Configuring the InnoDB Master Thread I/O Rate”.

- `innodb_flushing_avg_loops`

Command-Line Format	`--innodb-flushing-avg-loops=#`	
System Variable	Name	`innodb_flushing_avg_loops`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`30`
	Min Value	`1`
	Max Value	`1000`

Number of iterations for which `InnoDB` keeps the previously calculated snapshot of the flushing state, controlling how quickly adaptive flushing responds to changing workloads. Increasing the value makes the rate of flush operations change smoothly and gradually as the workload changes. Decreasing the value makes adaptive flushing adjust quickly to workload changes, which can cause spikes in flushing activity if the workload increases and decreases suddenly.

For related information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- `innodb_force_load_corrupted`

Command-Line Format	`--innodb-force-load-corrupted`	
System Variable	Name	`innodb_force_load_corrupted`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`OFF`

Permits `InnoDB` to load tables at startup that are marked as corrupted. Use only during troubleshooting, to recover data that is otherwise inaccessible. When troubleshooting is complete, disable this setting and restart the server.

- `innodb_force_recovery`

Command-Line Format	`--innodb-force-recovery=#`	
System Variable	Name	`innodb_force_recovery`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`0`

	Min Value	0
	Max Value	6

The crash recovery mode, typically only changed in serious troubleshooting situations. Possible values are from 0 to 6. For the meanings of these values and important information about `innodb_force_recovery`, see Section 15.20.2, “Forcing InnoDB Recovery”.

Warning

Only set this variable to a value greater than 0 in an emergency situation so that you can start `InnoDB` and dump your tables. As a safety measure, `InnoDB` prevents `INSERT`, `UPDATE`, or `DELETE` operations when `innodb_force_recovery` is greater than 0. An `innodb_force_recovery` setting of 4 or greater places `InnoDB` into read-only mode.

These restrictions may cause replication administration commands to fail with an error, as replication options such as `--relay-log-info-repository=TABLE` and `--master-info-repository=TABLE` store information in `InnoDB` tables.

- `innodb_ft_aux_table`

Command-Line Format	`--innodb-ft-aux-table=#`	
System Variable	**Name**	`innodb_ft_aux_table`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

Specifies the qualified name of an `InnoDB` table containing a `FULLTEXT` index. This variable is intended for diagnostic purposes.

After you set this variable to a name in the format *`db_name/table_name`*, the `INFORMATION_SCHEMA` tables `INNODB_FT_INDEX_TABLE`, `INNODB_FT_INDEX_CACHE`, `INNODB_FT_CONFIG`, `INNODB_FT_DELETED`, and `INNODB_FT_BEING_DELETED` show information about the search index for the specified table.

For more information, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

- `innodb_ft_cache_size`

Command-Line Format	`--innodb-ft-cache-size=#`	
System Variable	**Name**	`innodb_ft_cache_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	8000000

	Min Value	`1600000`
	Max Value	`80000000`

The memory allocated, in bytes, for the `InnoDB FULLTEXT` search index cache, which holds a parsed document in memory while creating an `InnoDB FULLTEXT` index. Index inserts and updates are only committed to disk when the `innodb_ft_cache_size` size limit is reached. `innodb_ft_cache_size` defines the cache size on a per table basis. To set a global limit for all tables, see `innodb_ft_total_cache_size`.

For more information, see InnoDB Full-Text Index Cache.

- `innodb_ft_enable_diag_print`

Command-Line Format	`--innodb-ft-enable-diag-print=#`	
System Variable	**Name**	`innodb_ft_enable_diag_print`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Whether to enable additional full-text search (FTS) diagnostic output. This option is primarily intended for advanced FTS debugging and will not be of interest to most users. Output is printed to the error log and includes information such as:

- FTS index sync progress (when the FTS cache limit is reached). For example:

```
FTS SYNC for table test, deleted count: 100 size: 10000 bytes
SYNC words: 100
```

- FTS optimize progress. For example:

```
FTS start optimize test
FTS_OPTIMIZE: optimize "mysql"
FTS_OPTIMIZE: processed "mysql"
```

- FTS index build progress. For example:

```
Number of doc processed: 1000
```

- For FTS queries, the query parsing tree, word weight, query processing time, and memory usage are printed. For example:

```
FTS Search Processing time: 1 secs: 100 millisec: row(s) 10000
Full Search Memory: 245666 (bytes),  Row: 10000
```

- `innodb_ft_enable_stopword`

Command-Line Format	`--innodb-ft-enable-stopword=#`	
System Variable	**Name**	`innodb_ft_enable_stopword`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Specifies that a set of stopwords is associated with an `InnoDB FULLTEXT` index at the time the index is created. If the `innodb_ft_user_stopword_table` option is set, the stopwords are taken from that table. Else, if the `innodb_ft_server_stopword_table` option is set, the stopwords are taken from that table. Otherwise, a built-in set of default stopwords is used.

For more information, see Section 12.9.4, “Full-Text Stopwords”.

- `innodb_ft_max_token_size`

Command-Line Format	`--innodb-ft-max-token-size=#`	
System Variable	**Name**	`innodb_ft_max_token_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`84`
	Min Value	`10`
	Max Value	`84`
Permitted Values	**Type**	integer
	Default	`84`
	Min Value	`10`
	Max Value	`84`

Maximum character length of words that are stored in an `InnoDB FULLTEXT` index. Setting a limit on this value reduces the size of the index, thus speeding up queries, by omitting long keywords or arbitrary collections of letters that are not real words and are not likely to be search terms.

For more information, see Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.

- `innodb_ft_min_token_size`

Command-Line Format	`--innodb-ft-min-token-size=#`	
System Variable	**Name**	`innodb_ft_min_token_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`3`
	Min Value	`0`

	Max Value	16

Minimum length of words that are stored in an InnoDB FULLTEXT index. Increasing this value reduces the size of the index, thus speeding up queries, by omitting common words that are unlikely to be significant in a search context, such as the English words “a” and “to”. For content using a CJK (Chinese, Japanese, Korean) character set, specify a value of 1.

For more information, see Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.

- innodb_ft_num_word_optimize

Command-Line Format	--innodb-ft-num-word-optimize=#	
System Variable	**Name**	innodb_ft_num_word_optimize
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	2000

Number of words to process during each OPTIMIZE TABLE operation on an InnoDB FULLTEXT index. Because a bulk insert or update operation to a table containing a full-text search index could require substantial index maintenance to incorporate all changes, you might do a series of OPTIMIZE TABLE statements, each picking up where the last left off.

For more information, see Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.

- innodb_ft_result_cache_limit

Command-Line Format	--innodb-ft-result-cache-limit=#	
System Variable	**Name**	innodb_ft_result_cache_limit
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	2000000000
	Min Value	1000000
	Max Value	2**32-1

The InnoDB full-text search query result cache limit (defined in bytes) per full-text search query or per thread. Intermediate and final InnoDB full-text search query results are handled in memory. Use innodb_ft_result_cache_limit to place a size limit on the full-text search query result cache to avoid excessive memory consumption in case of very large InnoDB full-text search query results (millions or hundreds of millions of rows, for example). Memory is allocated as required when a full-text search query is processed. If the result cache size limit is reached, an error is returned indicating that the query exceeds the maximum allowed memory.

The maximum value of innodb_ft_result_cache_limit for all platform types and bit sizes is 2**32-1.

- `innodb_ft_server_stopword_table`

Command-Line Format	`--innodb-ft-server-stopword-table=db_name/table_name`	
System Variable	Name	`innodb_ft_server_stopword_table`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string
	Default	`NULL`

This option is used to specify your own `InnoDB` `FULLTEXT` index stopword list for all `InnoDB` tables. To configure your own stopword list for a specific `InnoDB` table, use `innodb_ft_user_stopword_table`.

Set `innodb_ft_server_stopword_table` to the name of the table containing a list of stopwords, in the format *`db_name/table_name`*.

The stopword table must exist before you configure `innodb_ft_server_stopword_table`. `innodb_ft_enable_stopword` must be enabled and `innodb_ft_server_stopword_table` option must be configured before you create the `FULLTEXT` index.

The stopword table must be an `InnoDB` table, containing a single `VARCHAR` column named `value`.

For more information, see Section 12.9.4, “Full-Text Stopwords”.

- `innodb_ft_sort_pll_degree`

Command-Line Format	`--innodb-ft-sort-pll-degree=#`	
System Variable	Name	`innodb_ft_sort_pll_degree`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`2`
	Min Value	`1`
	Max Value	`32`

Number of threads used in parallel to index and tokenize text in an `InnoDB` `FULLTEXT` index when building a search index.

For related information, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and `innodb_sort_buffer_size`.

- `innodb_ft_total_cache_size`

Command-Line Format	`--innodb-ft-total-cache-size=#`	
System Variable	Name	`innodb_ft_total_cache_size`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`640000000`
	Min Value	`32000000`
	Max Value	`1600000000`

The total memory allocated, in bytes, for the `InnoDB` full-text search index cache for all tables. Creating numerous tables, each with a `FULLTEXT` search index, could consume a significant portion of available memory. `innodb_ft_total_cache_size` defines a global memory limit for all full-text search indexes to help avoid excessive memory consumption. If the global limit is reached by an index operation, a forced sync is triggered.

For more information, see InnoDB Full-Text Index Cache.

- `innodb_ft_user_stopword_table`

Command-Line Format	`--innodb-ft-user-stopword-table=db_name/table_name`	
System Variable	**Name**	`innodb_ft_user_stopword_table`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`NULL`

This option is used to specify your own `InnoDB` `FULLTEXT` index stopword list on a specific table. To configure your own stopword list for all `InnoDB` tables, use `innodb_ft_server_stopword_table`.

Set `innodb_ft_user_stopword_table` to the name of the table containing a list of stopwords, in the format *`db_name/table_name`*.

The stopword table must exist before you configure `innodb_ft_user_stopword_table`. `innodb_ft_enable_stopword` must be enabled and `innodb_ft_user_stopword_table` must be configured before you create the `FULLTEXT` index.

The stopword table must be an `InnoDB` table, containing a single `VARCHAR` column named `value`.

For more information, see Section 12.9.4, “Full-Text Stopwords”.

- `innodb_io_capacity`

Command-Line Format	`--innodb-io-capacity=#`	
System Variable	**Name**	`innodb_io_capacity`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`200`

	Min Value	`100`
	Max Value	`2**32-1`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`200`
	Min Value	`100`
	Max Value	`2**64-1`

The `innodb_io_capacity` parameter sets an upper limit on I/O activity performed by `InnoDB` background tasks, such as flushing pages from the buffer pool and merging data from the change buffer.

The `innodb_io_capacity` limit is a total limit for all buffer pool instances. When dirty pages are flushed, the limit is divided equally among buffer pool instances.

`innodb_io_capacity` should be set to approximately the number of I/O operations that the system can perform per second. Ideally, keep the setting as low as practical, but not so low that background activities fall behind. If the value is too high, data is removed from the buffer pool and insert buffer too quickly for caching to provide a significant benefit.

The default value is 200. For busy systems capable of higher I/O rates, you can set a higher value to help the server handle the background maintenance work associated with a high rate of row changes.

In general, you can increase the value as a function of the number of drives used for `InnoDB` I/O. For example, you can increase the value on systems that use multiple disks or solid-state disks (SSD).

The default setting of 200 is generally sufficient for a lower-end SSD. For a higher-end, bus-attached SSD, consider a higher setting such as 1000, for example. For systems with individual 5400 RPM or 7200 RPM drives, you might lower the value to `100`, which represents an estimated proportion of the I/O operations per second (IOPS) available to older-generation disk drives that can perform about 100 IOPS.

Although you can specify a very high value such as one million, in practice such large values have little if any benefit. Generally, a value of 20000 or higher is not recommended unless you have proven that lower values are insufficient for your workload.

Consider write workload when tuning `innodb_io_capacity`. Systems with large write workloads are likely to benefit from a higher setting. A lower setting may be sufficient for systems with a small write workload.

You can set `innodb_io_capacity` to any number 100 or greater to a maximum defined by `innodb_io_capacity_max`. `innodb_io_capacity` can be set in the MySQL option file (`my.cnf` or `my.ini`) or changed dynamically using a `SET GLOBAL` statement, which requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege.

The `innodb_flush_sync` configuration option causes the `innodb_io_capacity` setting to be ignored during bursts of I/O activity that occur at checkpoints. `innodb_flush_sync` is enabled by default.

See Section 15.6.8, “Configuring the InnoDB Master Thread I/O Rate” for more information. For general information about `InnoDB` I/O performance, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_io_capacity_max`

Command-Line Format	`--innodb-io-capacity-max=#`	
System Variable	**Name**	`innodb_io_capacity_max`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`see description`
	Min Value	`100`
	Max Value	`2**32-1`
Permitted Values (Unix, 64-bit platforms)	**Type**	integer
	Default	`see description`
	Min Value	`100`
	Max Value	`2**64-1`
Permitted Values (Windows, 64-bit platforms)	**Type**	integer
	Default	`see description`
	Min Value	`100`
	Max Value	`2**32-1`

If flushing activity falls behind, `InnoDB` can flush more aggressively than the limit imposed by `innodb_io_capacity`. `innodb_io_capacity_max` defines an upper limit for I/O capacity in such situations.

The `innodb_io_capacity_max` setting is a total limit for all buffer pool instances.

If you specify an `innodb_io_capacity` setting at startup but do not specify a value for `innodb_io_capacity_max`, `innodb_io_capacity_max` defaults to twice the value of `innodb_io_capacity`, with a minimum value of 2000.

When configuring `innodb_io_capacity_max`, twice the `innodb_io_capacity` is often a good starting point. The default value of 2000 is intended for workloads that use a solid-state disk (SSD) or more than one regular disk drive. A setting of 2000 is likely too high for workloads that do not use SSD or multiple disk drives, and could allow too much flushing. For a single regular disk drive, a setting between 200 and 400 is recommended. For a high-end, bus-attached SSD, consider a higher setting such as 2500. As with the `innodb_io_capacity` setting, keep the setting as low as practical, but not so low that `InnoDB` cannot sufficiently extend beyond the `innodb_io_capacity` limit, if necessary.

Consider write workload when tuning `innodb_io_capacity_max`. Systems with large write workloads may benefit from a higher setting. A lower setting may be sufficient for systems with a small write workload.

`innodb_io_capacity_max` cannot be set to a value lower than the `innodb_io_capacity` value.

Setting `innodb_io_capacity_max` to `DEFAULT` using a `SET` statement (`SET GLOBAL innodb_io_capacity_max=DEFAULT`) sets `innodb_io_capacity_max` to the maximum value.

For related information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- `innodb_limit_optimistic_insert_debug`

Command-Line Format	`--innodb-limit-optimistic-insert-debug=#`	
System Variable	Name	`innodb_limit_optimistic_insert_debug`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`2**32-1`

Limits the number of records per B-tree page. A default value of 0 means that no limit is imposed. This option is only available if debugging support is compiled in using the `WITH_DEBUG CMake` option.

- `innodb_lock_wait_timeout`

Command-Line Format	`--innodb-lock-wait-timeout=#`	
System Variable	Name	`innodb_lock_wait_timeout`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`50`
	Min Value	`1`
	Max Value	`1073741824`

The length of time in seconds an `InnoDB` transaction waits for a row lock before giving up. The default value is 50 seconds. A transaction that tries to access a row that is locked by another `InnoDB` transaction waits at most this many seconds for write access to the row before issuing the following error:

```
ERROR 1205 (HY000): Lock wait timeout exceeded; try restarting transaction
```

When a lock wait timeout occurs, the current statement is rolled back (not the entire transaction). To have the entire transaction roll back, start the server with the `--innodb_rollback_on_timeout` option. See also Section 15.20.4, “InnoDB Error Handling”.

You might decrease this value for highly interactive applications or OLTP systems, to display user feedback quickly or put the update into a queue for processing later. You might increase this value for long-running back-end operations, such as a transform step in a data warehouse that waits for other large insert or update operations to finish.

`innodb_lock_wait_timeout` applies to `InnoDB` row locks only. A MySQL table lock does not happen inside `InnoDB` and this timeout does not apply to waits for table locks.

The lock wait timeout value does not apply to deadlocks when `innodb_deadlock_detect` is enabled (the default) because `InnoDB` detects deadlocks immediately and rolls back one of the deadlocked transactions. When `innodb_deadlock_detect` is disabled, `InnoDB` relies on `innodb_lock_wait_timeout` for transaction rollback when a deadlock occurs. See Section 15.5.5.2, “Deadlock Detection and Rollback”.

`innodb_lock_wait_timeout` can be set at runtime with the `SET GLOBAL` or `SET SESSION` statement. Changing the `GLOBAL` setting requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege and affects the operation of all clients that subsequently connect. Any client can change the `SESSION` setting for `innodb_lock_wait_timeout`, which affects only that client.

- `innodb_log_buffer_size`

Command-Line Format	`--innodb-log-buffer-size=#`	
System Variable	Name	`innodb_log_buffer_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`16777216`
	Min Value	`1048576`
	Max Value	`4294967295`

The size in bytes of the buffer that `InnoDB` uses to write to the log files on disk. The default is 16MB. A large log buffer enables large transactions to run without the need to write the log to disk before the transactions commit. Thus, if you have transactions that update, insert, or delete many rows, making the log buffer larger saves disk I/O. For related information, see InnoDB Memory Configuration, and Section 8.5.4, “Optimizing InnoDB Redo Logging”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_log_checksums`

Command-Line Format	`--innodb-log-checksums=#`	
System Variable	Name	`innodb_log_checksums`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

Enables or disables checksums for redo log pages.

`innodb_log_checksums=ON` enables the `CRC-32C` checksum algorithm for redo log pages. When `innodb_log_checksums` is disabled, the contents of the redo log page checksum field are ignored.

Checksums on the redo log header page and redo log checkpoint pages are never disabled.

- `innodb_log_compressed_pages`

Command-Line Format	`--innodb-log-compressed-pages=#`	
System Variable	Name	`innodb_log_compressed_pages`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

Specifies whether images of re-compressed pages are written to the redo log. Re-compression may occur when changes are made to compressed data.

`innodb_log_compressed_pages` is enabled by default to prevent corruption that could occur if a different version of the `zlib` compression algorithm is used during recovery. If you are certain that the `zlib` version will not change, you can disable `innodb_log_compressed_pages` to reduce redo log generation for workloads that modify compressed data.

To measure the effect of enabling or disabling `innodb_log_compressed_pages`, compare redo log generation for both settings under the same workload. Options for measuring redo log generation include observing the `Log sequence number` (LSN) in the `LOG` section of `SHOW ENGINE INNODB STATUS` output, or monitoring `Innodb_os_log_written` status for the number of bytes written to the redo log files.

For related information, see Section 15.9.1.6, “Compression for OLTP Workloads”.

- `innodb_log_file_size`

Command-Line Format	`--innodb-log-file-size=#`	
System Variable	Name	`innodb_log_file_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`50331648`
	Min Value	`4194304`
	Max Value	`512GB / innodb_log_files_in_group`

The size in bytes of each log file in a log group. The combined size of log files (`innodb_log_file_size` * `innodb_log_files_in_group`) cannot exceed a maximum value that is slightly less than 512GB. A pair of 255 GB log files, for example, approaches the limit but does not exceed it. The default value is 48MB.

Generally, the combined size of the log files should be large enough that the server can smooth out peaks and troughs in workload activity, which often means that there is enough redo log space to handle more than an hour of write activity. The larger the value, the less checkpoint flush activity is required in the buffer pool, saving disk I/O. Larger log files also make crash recovery slower, although improvements to recovery performance in MySQL 5.5 and higher make the log file size less of a consideration.

The minimum `innodb_log_file_size` is 4MB.

For related information, see InnoDB Log File Configuration. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_log_files_in_group`

Command-Line Format	`--innodb-log-files-in-group=#`	
System Variable	Name	`innodb_log_files_in_group`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`2`
	Min Value	`2`
	Max Value	`100`

The number of log files in the log group. `InnoDB` writes to the files in a circular fashion. The default (and recommended) value is 2. The location of the files is specified by `innodb_log_group_home_dir`. The combined size of log files (`innodb_log_file_size` * `innodb_log_files_in_group`) can be up to 512GB.

For related information, see InnoDB Log File Configuration.

- `innodb_log_group_home_dir`

Command-Line Format	`--innodb-log-group-home-dir=dir_name`	
System Variable	Name	`innodb_log_group_home_dir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The directory path to the `InnoDB` redo log files, whose number is specified by `innodb_log_files_in_group`. If you do not specify any `InnoDB` log variables, the default is to create two files named `ib_logfile0` and `ib_logfile1` in the MySQL data directory. Log file size is given by the `innodb_log_file_size` system variable.

For related information, see InnoDB Log File Configuration.

- `innodb_log_write_ahead_size`

Command-Line Format	`--innodb-log-write-ahead-size=#`	
System Variable	Name	`innodb_log_write_ahead_size`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	integer
	Default	`8192`
	Min Value	`512 (log file block size)`
	Max Value	`Equal to innodb_page_size`

The write-ahead block size for the redo log, in bytes. To avoid “read-on-write”, set `innodb_log_write_ahead_size` to match the operating system or file system cache block size. Read-on-write occurs when redo log blocks are not entirely cached to the operating system or file system due to a mismatch between write-ahead block size for redo logs and operating system or file system cache block size.

Valid values for `innodb_log_write_ahead_size` are multiples of the `InnoDB` log file block size (2^n). The minimum value is the `InnoDB` log file block size (512). Write-ahead does not occur when the minimum value is specified. The maximum value is equal to `innodb_page_size`. If you specify a value for `innodb_log_write_ahead_size` that is larger than the `innodb_page_size` value, the `innodb_log_write_ahead_size` value is truncated to the `innodb_page_size` value.

Setting the `innodb_log_write_ahead_size` value too low in relation to the operating system or file system cache block size results in “read-on-write”. Setting the value too high may have a slight impact on `fsync` performance for log file writes due to several blocks being written at once.

- `innodb_lru_scan_depth`

Command-Line Format	`--innodb-lru-scan-depth=#`	
System Variable	**Name**	`innodb_lru_scan_depth`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`1024`
	Min Value	`100`
	Max Value	`2**32-1`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`1024`
	Min Value	`100`
	Max Value	`2**64-1`

A parameter that influences the algorithms and heuristics for the flush operation for the `InnoDB` buffer pool. Primarily of interest to performance experts tuning I/O-intensive workloads. It specifies, per buffer pool instance, how far down the buffer pool LRU list the page cleaner thread scans looking for dirty pages to flush. This is a background operation performed once per second.

A setting smaller than the default is generally suitable for most workloads. A value that is much higher than necessary may impact performance. Only consider increasing the value if you have spare I/O capacity under a typical workload. Conversely, if a write-intensive workload saturates your I/O capacity, decrease the value, especially in the case of a large buffer pool.

When tuning innodb_lru_scan_depth, start with a low value and configure the setting upward with the goal of rarely seeing zero free pages. Also, consider adjusting innodb_lru_scan_depth when changing the number of buffer pool instances, since innodb_lru_scan_depth * innodb_buffer_pool_instances defines the amount of work performed by the page cleaner thread each second.

For related information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- innodb_max_dirty_pages_pct

Command-Line Format	--innodb-max-dirty-pages-pct=#	
System Variable	Name	innodb_max_dirty_pages_pct
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	numeric
	Default	75
	Min Value	0
	Max Value	99.99

InnoDB tries to flush data from the buffer pool so that the percentage of dirty pages does not exceed this value. The default value is 75.

The innodb_max_dirty_pages_pct setting establishes a target for flushing activity. It does not affect the rate of flushing. For information about managing the rate of flushing, see Section 15.6.3.6, “Configuring InnoDB Buffer Pool Flushing”.

For related information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- innodb_max_dirty_pages_pct_lwm

Command-Line Format	--innodb-max-dirty-pages-pct-lwm=#	
System Variable	Name	innodb_max_dirty_pages_pct_lwm
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	numeric
	Default	0
	Min Value	0
	Max Value	99.99

Defines a low water mark representing the percentage of dirty pages at which preflushing is enabled to control the dirty page ratio. The default of 0 disables the pre-flushing behavior entirely. For more information, see Section 15.6.3.7, “Fine-tuning InnoDB Buffer Pool Flushing”.

- `innodb_max_purge_lag`

Command-Line Format	`--innodb-max-purge-lag=#`	
System Variable	Name	`innodb_max_purge_lag`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

Controls how to delay `INSERT`, `UPDATE`, and `DELETE` operations when purge operations are lagging (see Section 15.3, “InnoDB Multi-Versioning”). The default value is 0 (no delays).

The `InnoDB` transaction system maintains a list of transactions that have index records delete-marked by `UPDATE` or `DELETE` operations. The length of the list represents the *`purge_lag`* value. When *`purge_lag`* exceeds `innodb_max_purge_lag`, `INSERT`, `UPDATE`, and `DELETE` operations are delayed.

To prevent excessive delays in extreme situations where *`purge_lag`* becomes huge, you can limit the delay by setting the `innodb_max_purge_lag_delay` configuration option. The delay is computed at the beginning of a purge batch.

A typical setting for a problematic workload might be 1 million, assuming that transactions are small, only 100 bytes in size, and it is permissible to have 100MB of unpurged `InnoDB` table rows.

The lag value is displayed as the history list length in the `TRANSACTIONS` section of InnoDB Monitor output . For example, if the output includes the following lines, the lag value is 20:

```
------------
TRANSACTIONS
------------
Trx id counter 0 290328385
Purge done for trx's n:o < 0 290315608 undo n:o < 0 17
History list length 20
```

For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_max_purge_lag_delay`

Command-Line Format	`--innodb-max-purge-lag-delay=#`	
System Variable	Name	`innodb_max_purge_lag_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`

Specifies the maximum delay in milliseconds for the delay imposed by the `innodb_max_purge_lag` configuration option. A non-zero value represents an upper limit on the delay period computed from the formula based on the value of `innodb_max_purge_lag`. The default of zero means that there is no upper limit imposed on the delay interval.

For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_max_undo_log_size`

Command-Line Format	`--innodb-max-undo-log-size=#`	
System Variable	Name	`innodb_max_undo_log_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1073741824`
	Min Value	`10485760`
	Max Value	`2**64-1`

Defines a threshold size for undo tablespaces. If an undo tablespace exceeds the threshold, it can be marked for truncation when `innodb_undo_log_truncate` is enabled. The default value is 1024 MiB (1073741824 bytes).

For more information, see Section 15.7.8, “Truncating Undo Logs That Reside in Undo Tablespaces”.

- `innodb_merge_threshold_set_all_debug`

Command-Line Format	`--innodb-merge-threshold-set-all-debug=#`	
System Variable	Name	`innodb_merge_threshold_set_all_debug`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`50`
	Min Value	`1`
	Max Value	`50`

Overrides the current `MERGE_THRESHOLD` setting with the specified value for all indexes that are currently in the dictionary cache. This option is only available if debugging support is compiled in using the `WITH_DEBUG CMake` option. For related information, see Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

- `innodb_monitor_disable`

Command-Line Format	`--innodb-monitor-disable=[counter|module|pattern|all]`

System Variable	Name	innodb_monitor_disable
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

Disables InnoDB metrics counters. Counter data may be queried using the INFORMATION_SCHEMA.INNODB_METRICS table. For usage information, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

innodb_monitor_disable='latch' disables statistics collection for SHOW ENGINE INNODB MUTEX. For more information, see Section 13.7.5.15, “SHOW ENGINE Syntax”.

- innodb_monitor_enable

Command-Line Format	--innodb-monitor-enable=[counter|module|pattern|all]	
System Variable	Name	innodb_monitor_enable
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

Enables InnoDB metrics counters. Counter data may be queried using the INFORMATION_SCHEMA.INNODB_METRICS table. For usage information, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

innodb_monitor_enable='latch' enables statistics collection for SHOW ENGINE INNODB MUTEX. For more information, see Section 13.7.5.15, “SHOW ENGINE Syntax”.

- innodb_monitor_reset

Command-Line Format	--innodb-monitor-reset=[counter|module|pattern|all]	
System Variable	Name	innodb_monitor_reset
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

Resets the count value for InnoDB metrics counters to zero. Counter data may be queried using the INFORMATION_SCHEMA.INNODB_METRICS table. For usage information, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

innodb_monitor_reset='latch' resets statistics reported by SHOW ENGINE INNODB MUTEX. For more information, see Section 13.7.5.15, “SHOW ENGINE Syntax”.

- innodb_monitor_reset_all

Command-Line Format	--innodb-monitor-reset-all=[counter|module|pattern|all]	
System Variable	Name	innodb_monitor_reset_all

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

Resets all values (minimum, maximum, and so on) for `InnoDB` metrics counters. Counter data may be queried using the `INFORMATION_SCHEMA.INNODB_METRICS` table. For usage information, see Section 15.14.6, “InnoDB INFORMATION_SCHEMA Metrics Table”.

- `innodb_numa_interleave`

Command-Line Format	`--innodb-numa-interleave=#`	
System Variable	**Name**	`innodb_numa_interleave`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Enables the NUMA interleave memory policy for allocation of the `InnoDB` buffer pool. When `innodb_numa_interleave` is enabled, the NUMA memory policy is set to `MPOL_INTERLEAVE` for the `mysqld` process. After the `InnoDB` buffer pool is allocated, the NUMA memory policy is set back to `MPOL_DEFAULT`. For the `innodb_numa_interleave` option to be available, MySQL must be compiled on a NUMA-enabled Linux system.

`CMake` sets the default `WITH_NUMA` value based on whether the current platform has `NUMA` support. For more information, see Section 2.8.4, “MySQL Source-Configuration Options”.

- `innodb_old_blocks_pct`

Command-Line Format	`--innodb-old-blocks-pct=#`	
System Variable	**Name**	`innodb_old_blocks_pct`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`37`
	Min Value	`5`
	Max Value	`95`

Specifies the approximate percentage of the `InnoDB` buffer pool used for the old block sublist. The range of values is 5 to 95. The default value is 37 (that is, 3/8 of the pool). Often used in combination with `innodb_old_blocks_time`.

For more information, see Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”. For information about buffer pool management, the LRU algorithm, and eviction policies, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

- `innodb_old_blocks_time`

Command-Line Format	`--innodb-old-blocks-time=#`	
System Variable	Name	`innodb_old_blocks_time`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1000`
	Min Value	`0`
	Max Value	`2**32-1`

Non-zero values protect against the buffer pool being filled by data that is referenced only for a brief period, such as during a full table scan. Increasing this value offers more protection against full table scans interfering with data cached in the buffer pool.

Specifies how long in milliseconds a block inserted into the old sublist must stay there after its first access before it can be moved to the new sublist. If the value is 0, a block inserted into the old sublist moves immediately to the new sublist the first time it is accessed, no matter how soon after insertion the access occurs. If the value is greater than 0, blocks remain in the old sublist until an access occurs at least that many milliseconds after the first access. For example, a value of 1000 causes blocks to stay in the old sublist for 1 second after the first access before they become eligible to move to the new sublist.

The default value is 1000.

This configuration option is often used in combination with `innodb_old_blocks_pct`. For more information, see Section 15.6.3.4, “Making the Buffer Pool Scan Resistant”. For information about buffer pool management, the LRU algorithm, and eviction policies, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

- `innodb_online_alter_log_max_size`

Command-Line Format	`--innodb-online-alter-log-max-size=#`	
System Variable	Name	`innodb_online_alter_log_max_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`134217728`
	Min Value	`65536`
	Max Value	`2**64-1`

Specifies an upper limit on the size of the temporary log files used during online DDL operations for `InnoDB` tables. There is one such log file for each index being created or table being altered. This log file stores data inserted, updated, or deleted in the table during the DDL operation. The temporary log file is extended when needed by the value of `innodb_sort_buffer_size`, up

to the maximum specified by `innodb_online_alter_log_max_size`. If a temporary log file exceeds the upper size limit, the `ALTER TABLE` operation fails and all uncommitted concurrent DML operations are rolled back. Thus, a large value for this option allows more DML to happen during an online DDL operation, but also extends the period of time at the end of the DDL operation when the table is locked to apply the data from the log.

- `innodb_open_files`

Command-Line Format	`--innodb-open-files=#`	
System Variable	**Name**	`innodb_open_files`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`
	Min Value	`10`
	Max Value	`4294967295`

This configuration option is only relevant if you use multiple `InnoDB` tablespaces. It specifies the maximum number of `.ibd` files that MySQL can keep open at one time. The minimum value is 10. The default value is 300 if `innodb_file_per_table` is not enabled, and the higher of 300 and `table_open_cache` otherwise.

The file descriptors used for `.ibd` files are for `InnoDB` tables only. They are independent of those specified by the `--open-files-limit` server option, and do not affect the operation of the table cache. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_optimize_fulltext_only`

Command-Line Format	`--innodb-optimize-fulltext-only=#`	
System Variable	**Name**	`innodb_optimize_fulltext_only`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Changes the way `OPTIMIZE TABLE` operates on `InnoDB` tables. Intended to be enabled temporarily, during maintenance operations for `InnoDB` tables with `FULLTEXT` indexes.

By default, `OPTIMIZE TABLE` reorganizes data in the clustered index of the table. When this option is enabled, `OPTIMIZE TABLE` skips the reorganization of table data, and instead processes newly added, deleted, and updated token data for `InnoDB FULLTEXT` indexes. Fore more information, see Optimizing InnoDB Full-Text Indexes.

- `innodb_page_cleaners`

Command-Line Format	`--innodb-page-cleaners=#`	
System Variable	**Name**	`innodb_page_cleaners`

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`4`
	Min Value	`1`
	Max Value	`64`

The number of page cleaner threads that flush dirty pages from buffer pool instances. Page cleaner threads perform flush list and LRU flushing. When there are multiple page cleaner threads, buffer pool flushing tasks for each buffer pool instance are dispatched to idle page cleaner threads. The `innodb_page_cleaners` default value is 4. If the number of page cleaner threads exceeds the number of buffer pool instances, `innodb_page_cleaners` is automatically set to the same value as `innodb_buffer_pool_instances`.

If your workload is write-IO bound when flushing dirty pages from buffer pool instances to data files, and if your system hardware has available capacity, increasing the number of page cleaner threads may help improve write-IO throughput.

Multi-threaded page cleaner support extends to shutdown and recovery phases.

The `setpriority()` system call is used on Linux platforms where it is supported, and where the `mysqld` execution user is authorized, to give `page_cleaner` threads priority over other MySQL and InnoDB threads to help page flushing keep pace with the current workload. `mysqld` execution user authorization can be configured in `/etc/security/limits.conf`. For example, if `mysqld` is run under the `mysql` user, you can authorize the `mysql` user by adding lines similar to the following to `/etc/security/limits.conf`:

```
mysql              hard    nice       -20
mysql              soft    nice       -20
```

Refer to your Linux operating system documentation for more information.

- `innodb_page_size`

Command-Line Format	`--innodb-page-size=#k`	
System Variable	**Name**	`innodb_page_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	enumeration
	Default	`16384`
	Valid Values	`4k`
		`8k`
		`16k`
		`32k`
		`64k`
		`4096`

		8192
		16384
		32768
		65536

Specifies the page size for all InnoDB tablespaces in a MySQL instance. You can specify page size using the values 64k, 32k, 16k (the default), 8k, or 4k. Alternatively, you can specify page size in bytes (65536, 32768, 16384, 8192, 4096).

innodb_page_size can only be configured prior to initializing the MySQL instance and cannot be changed afterward. If no value is specified, the instance is initialized using the default page size. See Section 15.6.1, “InnoDB Startup Configuration”.

For both 32k and 64k page sizes, the maximum row length is approximately 16000 bytes. ROW_FORMAT=COMPRESSED is not supported when innodb_page_size is set to 32KB or 64KB. For innodb_page_size=32k, extent size is 2MB. For innodb_page_size=64k, extent size is 4MB. innodb_log_buffer_size should be set to at least 16M (the default) when using 32k or 64k page sizes.

The default 16KB page size or larger is appropriate for a wide range of workloads, particularly for queries involving table scans and DML operations involving bulk updates. Smaller page sizes might be more efficient for OLTP workloads involving many small writes, where contention can be an issue when single pages contain many rows. Smaller pages might also be efficient with SSD storage devices, which typically use small block sizes. Keeping the InnoDB page size close to the storage device block size minimizes the amount of unchanged data that is rewritten to disk.

The minimum file size for the first system tablespace data file (ibdata1) differs depending on the innodb_page_size value. See the innodb_data_file_path option description for more information.

For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- innodb_print_all_deadlocks

Command-Line Format	--innodb-print-all-deadlocks=#	
System Variable	**Name**	innodb_print_all_deadlocks
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	OFF

When this option is enabled, information about all deadlocks in InnoDB user transactions is recorded in the mysqld error log. Otherwise, you see information about only the last deadlock, using the SHOW ENGINE INNODB STATUS command. An occasional InnoDB deadlock is not necessarily an issue, because InnoDB detects the condition immediately and rolls back one of the transactions automatically. You might use this option to troubleshoot why deadlocks are occurring if an application does not have appropriate error-handling logic to detect the rollback and retry its operation. A large number of deadlocks might indicate the need to restructure transactions that issue DML or SELECT ... FOR UPDATE statements for multiple tables, so that each transaction accesses the tables in the same order, thus avoiding the deadlock condition.

For related information, see Section 15.5.5, “Deadlocks in InnoDB”.

- `innodb_purge_batch_size`

Command-Line Format	`--innodb-purge-batch-size=#`	
System Variable	Name	`innodb_purge_batch_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`300`
	Min Value	`1`
	Max Value	`5000`

Defines the number of undo log pages that purge parses and processes in one batch from the history list. In a multi-threaded purge configuration, the coordinator purge thread divides `innodb_purge_batch_size` by `innodb_purge_threads` and assigns that number of pages to each purge thread. The `innodb_purge_batch_size` option also defines the number of undo log pages that purge frees after every 128 iterations through the undo logs.

The `innodb_purge_batch_size` option is intended for advanced performance tuning in combination with the `innodb_purge_threads` setting. Most MySQL users need not change `innodb_purge_batch_size` from its default value.

For related information, see Section 15.6.10, “Configuring InnoDB Purge Scheduling”.

- `innodb_purge_threads`

Command-Line Format	`--innodb-purge-threads=#`	
System Variable	Name	`innodb_purge_threads`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`4`
	Min Value	`1`
	Max Value	`32`

The number of background threads devoted to the `InnoDB` purge operation. A minimum value of 1 signifies that the purge operation is always performed by a background thread, never as part of the master thread. Running the purge operation in one or more background threads helps reduce internal contention within `InnoDB`, improving scalability. Increasing the value to greater than 1 creates that many separate purge threads, which can improve efficiency on systems where DML operations are performed on multiple tables. The maximum is 32.

For related information, see Section 15.6.10, “Configuring InnoDB Purge Scheduling”.

- `innodb_purge_rseg_truncate_frequency`

Command-Line Format	--innodb-purge-rseg-truncate-frequency=#	
System Variable	Name	innodb_purge_rseg_truncate_frequency
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	128
	Min Value	1
	Max Value	128

Defines the frequency with which the purge system frees rollback segments. An undo tablespace cannot be truncated until its rollback segments are freed. Normally, the purge system frees rollback segments once every 128 times that purge is invoked. Reducing the `innodb_purge_rseg_truncate_frequency` value increases the frequency with which the purge thread frees rollback segments. The default value is 128.

`innodb_purge_rseg_truncate_frequency` is intended for use with `innodb_undo_log_truncate`. For more information, see Section 15.7.8, “Truncating Undo Logs That Reside in Undo Tablespaces”.

- `innodb_random_read_ahead`

Command-Line Format	--innodb-random-read-ahead=#	
System Variable	Name	innodb_random_read_ahead
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Enables the random read-ahead technique for optimizing `InnoDB` I/O.

For details about performance considerations for different types of read-ahead requests, see Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- `innodb_read_ahead_threshold`

Command-Line Format	--innodb-read-ahead-threshold=#	
System Variable	Name	innodb_read_ahead_threshold
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	56

	Min Value	0
	Max Value	64

Controls the sensitivity of linear read-ahead that InnoDB uses to prefetch pages into the buffer pool. If InnoDB reads at least innodb_read_ahead_threshold pages sequentially from an extent (64 pages), it initiates an asynchronous read for the entire following extent. The permissible range of values is 0 to 64. A value of 0 disables read-ahead. For the default of 56, InnoDB must read at least 56 pages sequentially from an extent to initiate an asynchronous read for the following extent.

Knowing how many pages are read through the read-ahead mechanism, and how many of these pages are evicted from the buffer pool without ever being accessed, can be useful when fine-tuning the innodb_read_ahead_threshold setting. SHOW ENGINE INNODB STATUS output displays counter information from the Innodb_buffer_pool_read_ahead and Innodb_buffer_pool_read_ahead_evicted global status variables, which report the number of pages brought into the buffer pool by read-ahead requests, and the number of such pages evicted from the buffer pool without ever being accessed, respectively. The status variables report global values since the last server restart.

SHOW ENGINE INNODB STATUS also shows the rate at which the read-ahead pages are read in and the rate at which such pages are evicted without being accessed. The per-second averages are based on the statistics collected since the last invocation of SHOW ENGINE INNODB STATUS and are displayed in the BUFFER POOL AND MEMORY section of the SHOW ENGINE INNODB STATUS output.

For more information, see Section 15.6.3.5, “Configuring InnoDB Buffer Pool Prefetching (Read-Ahead)”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

- innodb_read_io_threads

Command-Line Format	--innodb-read-io-threads=#	
System Variable	**Name**	innodb_read_io_threads
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	4
	Min Value	1
	Max Value	64

The number of I/O threads for read operations in InnoDB. Its counterpart for write threads is innodb_write_io_threads. For more information, see Section 15.6.6, “Configuring the Number of Background InnoDB I/O Threads”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

Note

On Linux systems, running multiple MySQL servers (typically more than 12) with default settings for innodb_read_io_threads, innodb_write_io_threads, and the Linux aio-max-nr setting can exceed system limits. Ideally, increase the aio-max-nr setting; as a

workaround, you might reduce the settings for one or both of the MySQL configuration options.

- `innodb_read_only`

Command-Line Format	`--innodb-read-only=#`	
System Variable	Name	`innodb_read_only`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`OFF`

Starts `InnoDB` in read-only mode. For distributing database applications or data sets on read-only media. Can also be used in data warehouses to share the same data directory between multiple instances. For more information, see Section 15.6.2, “Configuring InnoDB for Read-Only Operation”.

Previously, enabling the `innodb_read_only` system variable prevented creating and dropping tables only for the `InnoDB` storage engine. As of MySQL 8.0, enabling `innodb_read_only` prevents these operations for all storage engines. Table creation and drop operations for any storage engine modify data dictionary tables in the `mysql` system database, but those tables use the `InnoDB` storage engine and cannot be modified when `innodb_read_only` is enabled. The same principle applies to other table operations that require modifying data dictionary tables. Examples:

- If the `innodb_read_only` system variable is enabled, `ANALYZE TABLE` may fail because it cannot update statistics tables in the data dictionary, which use `InnoDB`. For `ANALYZE TABLE` operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a `MyISAM` table). To obtain the updated distribution statistics, set `information_schema_stats=LATEST`.
- `ALTER TABLE` *`tbl_name`* `ENGINE=`*`engine_name`* fails because it updates the storage engine designation, which is stored in the data dictionary.

In addition, other tables in the `mysql` system database use the `InnoDB` storage engine in MySQL 8.0. Making those tables read only results in restrictions on operations that modify them. Examples:

- Account-management statements such as `CREATE USER` and `GRANT` fail because the grant tables use `InnoDB`.
- The `INSTALL PLUGIN` and `UNINSTALL PLUGIN` plugin-management statements fail because the `plugin` table uses `InnoDB`.
- The `CREATE FUNCTION` and `DROP FUNCTION` UDF-management statements fail because the `func` table uses `InnoDB`.

- `innodb_redo_log_encrypt`

Introduced	8.0.1	
Command-Line Format	`--innodb-redo-log-encrypt=#`	
System Variable	Name	`innodb_redo_log_encrypt`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	boolean
	Default	`OFF`

Controls encryption of redo log data for tables encrypted using the `InnoDB` tablespace encryption feature. Encryption of redo log data is disabled by default. For more information, see Redo Log Data Encryption.

- `innodb_replication_delay`

Command-Line Format	`--innodb-replication-delay=#`	
System Variable	**Name**	`innodb_replication_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

The replication thread delay (in ms) on a slave server if `innodb_thread_concurrency` is reached.

- `innodb_rollback_on_timeout`

Command-Line Format	`--innodb-rollback-on-timeout`	
System Variable	**Name**	`innodb_rollback_on_timeout`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

`InnoDB` rolls back only the last statement on a transaction timeout by default. If `--innodb_rollback_on_timeout` is specified, a transaction timeout causes `InnoDB` to abort and roll back the entire transaction.

Note

If the start-transaction statement was `START TRANSACTION` or `BEGIN` statement, rollback does not cancel that statement. Further SQL statements become part of the transaction until the occurrence of `COMMIT`, `ROLLBACK`, or some SQL statement that causes an implicit commit.

For more information, see Section 15.20.4, “InnoDB Error Handling”.

- `innodb_rollback_segments`

Command-Line Format	`--innodb-rollback-segments=#`	
System Variable	**Name**	`innodb_rollback_segments`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`128`
	Min Value	`1`
	Max Value	`128`

As of MySQL 8.0.2, `innodb_rollback_segments` defines the number of rollback segments allocated to the system tablespace, the temporary tablespace, and each undo tablespace. Prior to MySQL 8.0.2, `innodb_rollback_segments` defines the total number of rollback segments used by `InnoDB`. Each rollback segment can support a maximum of 1023 data-modifying transactions.

Start with a low value for this parameter and gradually increase it to avoid allocating rollback segments that are not required. If `innodb_rollback_segments` is not set, it defaults to the maximum value of 128.

For more information about rollback segments, see Section 15.3, “InnoDB Multi-Versioning”. For information about configuring separate undo tablespaces, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

- `innodb_scan_directories`

Introduced	8.0.2	
Command-Line Format	`--innodb-scan-directories=#`	
System Variable	**Name**	`innodb_scan_directories`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`NULL`

If, during recovery, `InnoDB` encounters redo logs written since the last checkpoint, the redo logs must be applied to affected tablespaces. The process that identifies affected tablespaces is referred to as tablespace discovery. Tablespace discovery depends on tablespace map files that map tablespace IDs in the redo logs to tablespace files. If tablespace map files are lost or corrupted, the `innodb_scan_directories` startup option can be used to specify tablespace file directories. This option causes `InnoDB` to read the first page of each tablespace file in the specified directories and recreate tablespace map files so that the recovery process can apply redo logs to affected tablespaces.

`innodb_scan_directories` may be specified as an option in a startup command or in a MySQL configuration file. Quotes are used around the argument value because otherwise a semicolon (;) is interpreted as a special character by some command interpreters. (Unix shells treat it as a command terminator, for example.)

Startup command:

```
mysqld --innodb-scan-directories="directory_path_1;directory_path_2"
```

MySQL configuration file:

```
[mysqld]
innodb_scan_directories="directory_path_1;directory_path_2"
```

For more information, see Lost or Corrupted Tablespace Map Files.

- `innodb_saved_page_number_debug`

Command-Line Format	`--innodb-saved-page-number-debug=#`	
System Variable	Name	`innodb_saved_page_number_debug`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Max Value	`2**23-1`

Saves a page number. Setting the `innodb_fil_make_page_dirty_debug` option dirties the page defined by `innodb_saved_page_number_debug`. The `innodb_saved_page_number_debug` option is only available if debugging support is compiled in using the `WITH_DEBUG CMake` option.

- `innodb_sort_buffer_size`

Command-Line Format	`--innodb-sort-buffer-size=#`	
System Variable	Name	`innodb_sort_buffer_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`1048576`
	Min Value	`65536`
	Max Value	`67108864`

Specifies the size of sort buffers used to sort data during creation of an `InnoDB` index. The specified size defines the amount of data that is read into memory for internal sorting and then written out to disk. This process is referred to as a “run”. During the merge phase, pairs of buffers of the specified size are read in and merged. The larger the setting, the fewer runs and merges there are.

This sort area is only used for merge sorts during index creation, not during later index maintenance operations. Buffers are deallocated when index creation completes.

The value of this option also controls the amount by which the temporary log file is extended to record concurrent DML during online DDL operations.

Before this setting was made configurable, the size was hardcoded to 1048576 bytes (1MB), which remains the default.

During an `ALTER TABLE` or `CREATE TABLE` statement that creates an index, 3 buffers are allocated, each with a size defined by this option. Additionally, auxiliary pointers are allocated to rows

in the sort buffer so that the sort can run on pointers (as opposed to moving rows during the sort operation).

For a typical sort operation, a formula such as this one can be used to estimate memory consumption:

```
(6 /*FTS_NUM_AUX_INDEX*/ * (3*@@global.innodb_sort_buffer_size)
+ 2 * number_of_partitions * number_of_secondary_indexes_created
* (@@global.innodb_sort_buffer_size/dict_index_get_min_size(index)*/)
* 8 /*64-bit sizeof *buf->tuples*/")
```

`@@global.innodb_sort_buffer_size/dict_index_get_min_size(index)` indicates the maximum tuples held. `2 * (@@global.innodb_sort_buffer_size/ *dict_index_get_min_size(index)*/) * 8 /*64-bit size of *buf->tuples*/` indicates auxiliary pointers allocated.

Note

For 32-bit, multiply by 4 instead of 8.

For parallel sorts on a full-text index, multiply by the `innodb_ft_sort_pll_degree` setting:

```
(6 /*FTS_NUM_AUX_INDEX*/ * @@global.innodb_ft_sort_pll_degree)
```

- `innodb_spin_wait_delay`

Command-Line Format	`--innodb-spin-wait-delay=#`	
System Variable	**Name**	`innodb_spin_wait_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`6`
	Min Value	`0`
	Max Value	`2**32-1`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`6`
	Min Value	`0`
	Max Value	`2**64-1`

The maximum delay between polls for a spin lock. The low-level implementation of this mechanism varies depending on the combination of hardware and operating system, so the delay does not correspond to a fixed time interval. For more information, see Section 15.6.9, “Configuring Spin Lock Polling”.

- `innodb_stats_auto_recalc`

Command-Line Format	`--innodb-stats-auto-recalc=#`	
System Variable	**Name**	`innodb_stats_auto_recalc`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Causes `InnoDB` to automatically recalculate persistent statistics after the data in a table is changed substantially. The threshold value is 10% of the rows in the table. This setting applies to tables created when the `innodb_stats_persistent` option is enabled. Automatic statistics recalculation may also be configured by specifying `STATS_PERSISTENT=1` in a `CREATE TABLE` or `ALTER TABLE` statement. The amount of data sampled to produce the statistics is controlled by the `innodb_stats_persistent_sample_pages` configuration option.

For more information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- `innodb_stats_include_delete_marked`

Introduced	8.0.1	
Command-Line Format	`--innodb-stats-include-delete-marked=#`	
System Variable	**Name**	`innodb_stats_include_delete_marked`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

By default, `InnoDB` reads uncommitted data when calculating statistics. In the case of an uncommitted transaction that deletes rows from a table, `InnoDB` excludes records that are delete-marked when calculating row estimates and index statistics, which can lead to non-optimal execution plans for other transactions that are operating on the table concurrently using a transaction isolation level other than `READ UNCOMMITTED`. To avoid this scenario, `innodb_stats_include_delete_marked` can be enabled to ensure that `InnoDB` includes delete-marked records when calculating persistent optimizer statistics.

When `innodb_stats_include_delete_marked` is enabled, `ANALYZE TABLE` considers delete-marked records when recalculating statistics.

`innodb_stats_include_delete_marked` is a global setting that affects all `InnoDB` tables. It is only applicable to persistent optimizer statistics.

For related information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- `innodb_stats_method`

Command-Line Format	`--innodb-stats-method=name`	
System Variable	**Name**	`innodb_stats_method`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	enumeration
	Default	`nulls_equal`
	Valid Values	`nulls_equal`
		`nulls_unequal`
		`nulls_ignored`

How the server treats `NULL` values when collecting statistics about the distribution of index values for `InnoDB` tables. Permitted values are `nulls_equal`, `nulls_unequal`, and `nulls_ignored`. For `nulls_equal`, all `NULL` index values are considered equal and form a single value group with a size equal to the number of `NULL` values. For `nulls_unequal`, `NULL` values are considered unequal, and each `NULL` forms a distinct value group of size 1. For `nulls_ignored`, `NULL` values are ignored.

The method used to generate table statistics influences how the optimizer chooses indexes for query execution, as described in Section 8.3.7, “InnoDB and MyISAM Index Statistics Collection”.

- `innodb_stats_on_metadata`

Command-Line Format	`--innodb-stats-on-metadata`	
System Variable	**Name**	`innodb_stats_on_metadata`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

This option only applies when optimizer statistics are configured to be non-persistent. Optimizer statistics are not persisted to disk when `innodb_stats_persistent` is disabled or when individual tables are created or altered with `STATS_PERSISTENT=0`. For more information, see Section 15.6.11.2, “Configuring Non-Persistent Optimizer Statistics Parameters”.

When `innodb_stats_on_metadata` is enabled, `InnoDB` updates non-persistent statistics when metadata statements such as `SHOW TABLE STATUS` or when accessing the `INFORMATION_SCHEMA.TABLES` or `INFORMATION_SCHEMA.STATISTICS` tables. (These updates are similar to what happens for `ANALYZE TABLE`.) When disabled, `InnoDB` does not update statistics during these operations. Leaving the setting disabled can improve access speed for schemas that have a large number of tables or indexes. It can also improve the stability of execution plans for queries that involve `InnoDB` tables.

To change the setting, issue the statement `SET GLOBAL innodb_stats_on_metadata=`*`mode`*, where *`mode`* is either `ON` or `OFF` (or `1` or `0`). Changing the setting requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege and immediately affects the operation of all connections.

- `innodb_stats_persistent`

Command-Line Format	`--innodb-stats-persistent=setting`	
System Variable	**Name**	`innodb_stats_persistent`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	boolean
	Default	`ON`
	Valid Values	`OFF`
		`ON`
		`0`
		`1`

Specifies whether `InnoDB` index statistics are persisted to disk. Otherwise, statistics may be recalculated frequently which can lead to variations in query execution plans. This setting is stored with each table when the table is created. You can set `innodb_stats_persistent` at the global level before creating a table, or use the `STATS_PERSISTENT` clause of the `CREATE TABLE` and `ALTER TABLE` statements to override the system-wide setting and configure persistent statistics for individual tables.

For more information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

- `innodb_stats_persistent_sample_pages`

Command-Line Format	`--innodb-stats-persistent-sample-pages=#`	
System Variable	**Name**	`innodb_stats_persistent_sample_pages`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`20`

The number of index pages to sample when estimating cardinality and other statistics for an indexed column, such as those calculated by `ANALYZE TABLE`. Increasing the value improves the accuracy of index statistics, which can improve the query execution plan, at the expense of increased I/O during the execution of `ANALYZE TABLE` for an `InnoDB` table. For more information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.

Note

Setting a high value for `innodb_stats_persistent_sample_pages` could result in lengthy `ANALYZE TABLE` execution time. To estimate the number of database pages accessed by `ANALYZE TABLE`, see Section 15.6.11.3, “Estimating ANALYZE TABLE Complexity for InnoDB Tables”.

`innodb_stats_persistent_sample_pages` only applies when `innodb_stats_persistent` is enabled for a table; when `innodb_stats_persistent` is disabled, `innodb_stats_transient_sample_pages` applies instead.

- `innodb_stats_transient_sample_pages`

Command-Line Format	`--innodb-stats-transient-sample-pages=#`	
System Variable	**Name**	`innodb_stats_transient_sample_pages`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	integer
	Default	8

The number of index pages to sample when estimating cardinality and other statistics for an indexed column, such as those calculated by `ANALYZE TABLE`. The default value is 8. Increasing the value improves the accuracy of index statistics, which can improve the query execution plan, at the expense of increased I/O when opening an `InnoDB` table or recalculating statistics. For more information, see Section 15.6.11.2, “Configuring Non-Persistent Optimizer Statistics Parameters”.

Note

Setting a high value for `innodb_stats_transient_sample_pages` could result in lengthy `ANALYZE TABLE` execution time. To estimate the number of database pages accessed by `ANALYZE TABLE`, see Section 15.6.11.3, “Estimating ANALYZE TABLE Complexity for InnoDB Tables”.

`innodb_stats_transient_sample_pages` only applies when `innodb_stats_persistent` is disabled for a table; when `innodb_stats_persistent` is enabled, `innodb_stats_persistent_sample_pages` applies instead. Takes the place of `innodb_stats_sample_pages`. For more information, see Section 15.6.11.2, “Configuring Non-Persistent Optimizer Statistics Parameters”.

- `innodb_status_output`

Command-Line Format	`--innodb-status-output`	
System Variable	**Name**	`innodb_status_output`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Enables or disables periodic output for the standard `InnoDB` Monitor. Also used in combination with `innodb_status_output_locks` to enable or disable periodic output for the `InnoDB` Lock Monitor. For more information, see Section 15.16.2, “Enabling InnoDB Monitors”.

- `innodb_status_output_locks`

Command-Line Format	`--innodb-status-output-locks`	
System Variable	**Name**	`innodb_status_output_locks`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Enables or disables the `InnoDB` Lock Monitor. When enabled, the `InnoDB` Lock Monitor prints additional information about locks in `SHOW ENGINE INNODB STATUS` output and in periodic output printed to the MySQL error log. Periodic output for the `InnoDB` Lock Monitor is printed as part of the standard `InnoDB` Monitor output. The standard `InnoDB` Monitor must therefore be enabled for the `InnoDB` Lock Monitor to print data to the MySQL error log periodically. For more information, see Section 15.16.2, “Enabling InnoDB Monitors”.

- `innodb_strict_mode`

Command-Line Format	`--innodb-strict-mode=#`	
System Variable	Name	`innodb_strict_mode`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`ON`

When `innodb_strict_mode` is enabled, `InnoDB` returns errors rather than warnings for certain conditions.

Strict mode helps guard against ignored typos and syntax errors in SQL, or other unintended consequences of various combinations of operational modes and SQL statements. When `innodb_strict_mode` is enabled, `InnoDB` raises error conditions in certain cases, rather than issuing a warning and processing the specified statement (perhaps with unintended behavior). This is analogous to `sql_mode` in MySQL, which controls what SQL syntax MySQL accepts, and determines whether it silently ignores errors, or validates input syntax and data values.

The `innodb_strict_mode` setting affects the handling of syntax errors for `CREATE TABLE`, `ALTER TABLE`, `CREATE INDEX`, and `OPTIMIZE TABLE` statements. `innodb_strict_mode` also enables a record size check, so that an `INSERT` or `UPDATE` never fails due to the record being too large for the selected page size.

Oracle recommends enabling `innodb_strict_mode` when using `ROW_FORMAT` and `KEY_BLOCK_SIZE` clauses in `CREATE TABLE`, `ALTER TABLE`, and `CREATE INDEX` statements. When `innodb_strict_mode` is disabled, `InnoDB` ignores conflicting clauses and creates the table or index with only a warning in the message log. The resulting table might have different characteristics than intended, such as lack of compression support when attempting to create a compressed table. When `innodb_strict_mode` is enabled, such problems generate an immediate error and the table or index is not created.

You can enable or disable `innodb_strict_mode` on the command line when starting `mysqld`, or in a MySQL configuration file. You can also enable or disable `innodb_strict_mode` at runtime with the statement `SET [GLOBAL|SESSION] innodb_strict_mode=`*`mode`*, where *`mode`* is either `ON` or `OFF`. Changing the `GLOBAL` setting requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege and affects the operation of all clients that subsequently connect. Any client can change the `SESSION` setting for `innodb_strict_mode`, and the setting affects only that client.

`innodb_strict_mode` is not applicable to general tablespaces. Tablespace management rules for general tablespaces are strictly enforced independently of `innodb_strict_mode`. For more information, see Section 13.1.16, “CREATE TABLESPACE Syntax”.

- `innodb_sync_array_size`

Command-Line Format	`--innodb-sync-array-size=#`	
System Variable	Name	`innodb_sync_array_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`1`

	Min Value	1
	Max Value	1024

Splits an internal data structure used to coordinate threads, for higher concurrency in workloads with large numbers of waiting threads. This setting must be configured when the MySQL instance is starting up, and cannot be changed afterward. Increasing the value is recommended for workloads that frequently produce a large number of waiting threads, typically greater than 768.

- `innodb_sync_spin_loops`

Command-Line Format	`--innodb-sync-spin-loops=#`	
System Variable	**Name**	`innodb_sync_spin_loops`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`30`
	Min Value	`0`
	Max Value	`4294967295`

The number of times a thread waits for an `InnoDB` mutex to be freed before the thread is suspended.

- `innodb_sync_debug`

Command-Line Format	`--innodb-sync-debug=#`	
System Variable	**Name**	`innodb_sync_debug`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Enables sync debug checking for the `InnoDB` storage engine. This option is only available if debugging support is compiled in using the `WITH_DEBUG` `CMake` option.

- `innodb_table_locks`

Command-Line Format	`--innodb-table-locks`	
System Variable	**Name**	`innodb_table_locks`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean

	Default	TRUE

If autocommit = 0, InnoDB honors LOCK TABLES; MySQL does not return from LOCK TABLES ... WRITE until all other threads have released all their locks to the table. The default value of innodb_table_locks is 1, which means that LOCK TABLES causes InnoDB to lock a table internally if autocommit = 0.

In MySQL 8.0, innodb_table_locks = 0 has no effect for tables locked explicitly with LOCK TABLES ... WRITE. It does have an effect for tables locked for read or write by LOCK TABLES ... WRITE implicitly (for example, through triggers) or by LOCK TABLES ... READ.

For related information, see Section 15.5, “InnoDB Locking and Transaction Model”.

- innodb_temp_data_file_path

Command-Line Format	--innodb-temp-data-file-path=file	
System Variable	**Name**	innodb_temp_data_file_path
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	ibtmp1:12M:autoextend

Specifies the path, file name, and file size for InnoDB temporary tablespace data files. The full directory path for a file is formed by concatenating innodb_data_home_dir to the path specified by innodb_temp_data_file_path. File size is specified in KB, MB, or GB (1024MB) by appending K, M, or G to the size value. The sum of the sizes of the files must be slightly larger than 12MB. If you do not specify a value for innodb_temp_data_file_path, the default behavior is to create a single auto-extending temporary tablespace data file, slightly larger than 12MB, named ibtmp1. The size limit of individual files is determined by your operating system. You can set the file size to more than 4GB on operating systems that support big files. Use of raw disk partitions for temporary tablespace data files is not supported.

A temporary tablespace data file name cannot be the same as an InnoDB data file name. Any inability or error creating a temporary tablespace data file is treated as fatal and server startup is refused. The temporary tablespace has a dynamically generated space ID, which can change on each server restart.

All InnoDB temporary tables are created in the shared temporary tablespace. InnoDB does not support compressed temporary tables.

Metadata about active InnoDB temporary tables is located in INFORMATION_SCHEMA.INNODB_TEMP_TABLE_INFO.

For related information, see Section 15.4.11, “Temporary Tablespace”.

- innodb_tmpdir

Command-Line Format	--innodb-tmpdir=path	
System Variable	**Name**	innodb_tmpdir
	Variable Scope	Global, Session
	Dynamic Variable	Yes

Permitted Values	Type	directory name
	Default	NULL

Used to define an alternate directory for temporary sort files created during online ALTER TABLE operations that rebuild the table.

Online ALTER TABLE operations that rebuild the table also create an *intermediate* table file in the same directory as the original table. The innodb_tmpdir option is not applicable to intermediate table files.

A valid value is any directory path other than the MySQL data directory path. If the value is NULL (the default), temporary files are created MySQL temporary directory ($TMPDIR on Unix, %TEMP% on Windows, or the directory specified by the --tmpdir configuration option). If a directory is specified, existence of the directory and permissions are only checked when innodb_tmpdir is configured using a SET statement. If a symlink is provided in a directory string, the symlink is resolved and stored as an absolute path. The path should not exceed 512 bytes. An online ALTER TABLE operation reports an error if innodb_tmpdir is set to an invalid directory. innodb_tmpdir overrides the MySQL tmpdir setting but only for online ALTER TABLE operations.

The FILE privilege is required to configure innodb_tmpdir.

The innodb_tmpdir option was introduced to help avoid overflowing a temporary file directory located on a tmpfs file system. Such overflows could occur as a result of large temporary sort files created during online ALTER TABLE operations that rebuild the table.

In replication environments, only consider replicating the innodb_tmpdir setting if all servers have the same operating system environment. Otherwise, replicating the innodb_tmpdir setting could result in a replication failure when running online ALTER TABLE operations that rebuild the table. If server operating environments differ, it is recommended that you configure innodb_tmpdir on each server individually.

For more information, see Where InnoDB Stores Temporary Files. For information about online ALTER TABLE operations, see Section 15.12, “InnoDB and Online DDL”.

- innodb_thread_concurrency

Command-Line Format	--innodb-thread-concurrency=#	
System Variable	Name	innodb_thread_concurrency
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	0
	Min Value	0
	Max Value	1000

InnoDB tries to keep the number of operating system threads concurrently inside InnoDB less than or equal to the limit given by this variable (InnoDB uses operating system threads to process user transactions). Once the number of threads reaches this limit, additional threads are placed into a wait state within a “First In, First Out” (FIFO) queue for execution. Threads waiting for locks are not counted in the number of concurrently executing threads.

The range of this variable is 0 to 1000. A value of 0 (the default) is interpreted as infinite concurrency (no concurrency checking). Disabling thread concurrency checking enables `InnoDB` to create as many threads as it needs. A value of 0 also disables the `queries inside InnoDB` and `queries in queue counters` in the `ROW OPERATIONS` section of `SHOW ENGINE INNODB STATUS` output.

Consider setting this variable if your MySQL instance shares CPU resources with other applications, or if your workload or number of concurrent users is growing. The correct setting depends on workload, computing environment, and the version of MySQL that you are running. You will need to test a range of values to determine the setting that provides the best performance. `innodb_thread_concurrency` is a dynamic variable, which allows you to experiment with different settings on a live test system. If a particular setting performs poorly, you can quickly set `innodb_thread_concurrency` back to 0.

Use the following guidelines to help find and maintain an appropriate setting:

- If the number of concurrent user threads for a workload is less than 64, set `innodb_thread_concurrency=0`.
- If your workload is consistently heavy or occasionally spikes, start by setting `innodb_thread_concurrency=128` and then lowering the value to 96, 80, 64, and so on, until you find the number of threads that provides the best performance. For example, suppose your system typically has 40 to 50 users, but periodically the number increases to 60, 70, or even 200. You find that performance is stable at 80 concurrent users but starts to show a regression above this number. In this case, you would set `innodb_thread_concurrency=80` to avoid impacting performance.
- If you do not want `InnoDB` to use more than a certain number of vCPUs for user threads (20 vCPUs, for example), set `innodb_thread_concurrency` to this number (or possibly lower, depending on performance results). If your goal is to isolate MySQL from other applications, you may consider binding the `mysqld` process exclusively to the vCPUs. Be aware, however, that exclusive binding could result in non-optimal hardware usage if the `mysqld` process is not consistently busy. In this case, you might bind the `mysqld` process to the vCPUs but also allow other applications to use some or all of the vCPUs.

  Note

  From an operating system perspective, using a resource management solution to manage how CPU time is shared among applications may be preferable to binding the `mysqld` process. For example, you could assign 90% of vCPU time to a given application while other critical process *are not* running, and scale that value back to 40% when other critical processes *are* running.

- `innodb_thread_concurrency` values that are too high can cause performance regression due to increased contention on system internals and resources.
- In some cases, the optimal `innodb_thread_concurrency` setting can be smaller than the number of vCPUs.
- Monitor and analyze your system regularly. Changes to workload, number of users, or computing environment may require that you adjust the `innodb_thread_concurrency` setting.

For related information, see Section 15.6.5, “Configuring Thread Concurrency for InnoDB”.

- `innodb_trx_purge_view_update_only_debug`

Command-Line Format	`--innodb-trx-purge-view-update-only-debug=#`

System Variable	Name	innodb_trx_purge_view_update_only_debug
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Pauses purging of delete-marked records while allowing the purge view to be updated. This option artificially creates a situation in which the purge view is updated but purges have not yet been performed. This option is only available if debugging support is compiled in using the WITH_DEBUG CMake option.

- innodb_trx_rseg_n_slots_debug

Command-Line Format	--innodb-trx-rseg-n-slots-debug=#	
System Variable	Name	innodb_trx_rseg_n_slots_debug
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	0
	Max Value	1024

Sets a debug flag that limits TRX_RSEG_N_SLOTS to a given value for the trx_rsegf_undo_find_free function that looks for free slots for undo log segments. This option is only available if debugging support is compiled in using the WITH_DEBUG CMake option.

- innodb_thread_sleep_delay

Command-Line Format	--innodb-thread-sleep-delay=#	
System Variable	Name	innodb_thread_sleep_delay
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	10000
	Min Value	0
	Max Value	1000000

How long InnoDB threads sleep before joining the InnoDB queue, in microseconds. The default value is 10000. A value of 0 disables sleep. You can set the configuration option innodb_adaptive_max_sleep_delay to the highest value you would allow for innodb_thread_sleep_delay, and InnoDB automatically adjusts innodb_thread_sleep_delay up or down depending on current thread-scheduling activity. This dynamic adjustment helps the thread scheduling mechanism to work smoothly during times when the system is lightly loaded or when it is operating near full capacity.

For more information, see Section 15.6.5, “Configuring Thread Concurrency for InnoDB”.

- `innodb_undo_directory`

Command-Line Format	`--innodb-undo-directory=dir_name`	
System Variable	Name	`innodb_undo_directory`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	directory name

The path where `InnoDB` creates undo tablespaces. Typically used to place undo logs on a different storage device. Used in conjunction with `innodb_rollback_segments` and `innodb_undo_tablespaces`, which determine the disk layout of the undo logs outside the system tablespace.

There is no default value (it is NULL). If a path is not specified, undo tablespaces are created in the MySQL data directory, as defined by `datadir`.

For more information, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

- `innodb_undo_log_encrypt`

Introduced	8.0.1	
Command-Line Format	`--innodb-undo-log-encrypt=#`	
System Variable	Name	`innodb_undo_log_encrypt`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Controls encryption of undo log data for tables encrypted using the `InnoDB` tablespace encryption feature. Only applies to undo logs that reside in separate undo tablespaces. See Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”. Encryption is not supported for undo log data that resides in the system tablespace. For more information, see Undo Log Data Encryption.

- `innodb_undo_log_truncate`

Command-Line Format	`--innodb-undo-log-truncate=#`	
System Variable	Name	`innodb_undo_log_truncate`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (<= 8.0.1)	Type	boolean
	Default	`OFF`
Permitted Values (>= 8.0.2)	Type	boolean
	Default	`ON`

When enabled, undo tablespaces that exceed the threshold value defined by `innodb_max_undo_log_size` are marked for truncation. Only undo logs that reside in undo tablespaces can be truncated. Truncating undo logs that reside in the system tablespace is not supported. For truncation to occur, there must be at least two undo tablespaces.

The `innodb_purge_rseg_truncate_frequency` configuration option can be used to expedite truncation of undo tablepaces.

For more information, see Section 15.7.8, “Truncating Undo Logs That Reside in Undo Tablespaces”.

- `innodb_undo_logs`

Deprecated	5.7.19	
Removed	8.0.2	
Command-Line Format	`--innodb-undo-logs=#`	
System Variable	**Name**	`innodb_undo_logs`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`128`
	Min Value	`0`
	Max Value	`128`

Note

`innodb_undo_logs` is removed in MySQL 8.0.2.

Defines the number of rollback segments used by `InnoDB`. The `innodb_undo_logs` option is an alias for `innodb_rollback_segments`. For more information, see the description of `innodb_rollback_segments`.

- `innodb_undo_tablespaces`

Command-Line Format	`--innodb-undo-tablespaces=#`	
System Variable (<= 8.0.1)	**Name**	`innodb_undo_tablespaces`
	Variable Scope	Global
	Dynamic Variable	No
System Variable (>= 8.0.2)	**Name**	`innodb_undo_tablespaces`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (<= 8.0.1)	**Type**	integer
	Default	`0`

	Min Value	0
	Max Value	95
Permitted Values (>= 8.0.2)	**Type**	integer
	Default	2
	Min Value	0
	Max Value	127

The number of undo tablespaces used by InnoDB. The default value is 2 as of MySQL 8.0.2, and 0 before that.

Undo logs can become large during long-running transactions. Using multiple undo tablespaces reduces the size of any one undo tablespace.

Undo tablespace files are created in the location defined by innodb_undo_directory. File names are in the form of undo*NNN*, where *NNN* is the undo tablespace number. The default size of an undo tablespace file is 10MiB.

As of MySQL 8.0.2, you can configure innodb_undo_tablespaces at startup or while the server is running. Increasing the innodb_undo_tablespaces setting creates the specified number of undo tablespaces and adds them to the list of active undo tablespaces. Decreasing the innodb_undo_tablespaces setting removes undo tablespaces from the list of active undo tablespaces. However, these undo tablespaces remain active until they are no longer used by existing transactions. Undo tablespaces are made inactive rather than deleted so that the number of active undo tablespaces can be increased again easily.

Prior to MySQL 8.0.2, you can only configure innodb_undo_tablespaces when initializing the MySQL instance. The setting cannot be changed afterward. If no value is specified, the instance is initialized using the default setting of 0. Attempting to restart InnoDB with a greater number of undo tablespaces than specified when the MySQL instance was initialized results in a startup failure and an error stating that InnoDB did not find the expected number of undo tablespaces.

If innodb_undo_tablespaces is set to 0, the system tablespace is used for rollback segments. A setting greater than 0 creates that number of undo tablespaces. A value greater than 0 also means that rollback segments in the system tablespace are no longer assigned to transactions.

Note

Setting innodb_undo_tablespaces to 0 is deprecated and will not be supported in a future release.

For more information, see Section 15.7.7, “Storing InnoDB Undo Logs in Separate Tablespaces”.

- innodb_use_native_aio

Command-Line Format	--innodb-use-native-aio=#	
System Variable	**Name**	innodb_use_native_aio
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean

	Default	ON

Specifies whether to use the Linux asynchronous I/O subsystem. This variable applies to Linux systems only, and cannot be changed while the server is running. Normally, you do not need to configure this option, because it is enabled by default.

The asynchronous I/O capability that InnoDB has on Windows systems is available on Linux systems. (Other Unix-like systems continue to use synchronous I/O calls.) This feature improves the scalability of heavily I/O-bound systems, which typically show many pending reads/writes in SHOW ENGINE INNODB STATUS\G output.

Running with a large number of InnoDB I/O threads, and especially running multiple such instances on the same server machine, can exceed capacity limits on Linux systems. In this case, you may receive the following error:

```
EAGAIN: The specified maxevents exceeds the user's limit of available events.
```

You can typically address this error by writing a higher limit to /proc/sys/fs/aio-max-nr.

However, if a problem with the asynchronous I/O subsystem in the OS prevents InnoDB from starting, you can start the server with innodb_use_native_aio=0. This option may also be disabled automatically during startup if InnoDB detects a potential problem such as a combination of tmpdir location, tmpfs file system, and Linux kernel that does not support AIO on tmpfs.

For more information, see Section 15.6.7, “Using Asynchronous I/O on Linux”.

- innodb_version

  The InnoDB version number. In MySQL 8.0, separate version numbering for InnoDB does not apply and this value is the same the version number of the server.

- innodb_write_io_threads

Command-Line Format	--innodb-write-io-threads=#	
System Variable	**Name**	innodb_write_io_threads
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	4
	Min Value	1
	Max Value	64

The number of I/O threads for write operations in InnoDB. The default value is 4. Its counterpart for read threads is innodb_read_io_threads. For more information, see Section 15.6.6, “Configuring the Number of Background InnoDB I/O Threads”. For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

Note

On Linux systems, running multiple MySQL servers (typically more than 12) with default settings for innodb_read_io_threads, innodb_write_io_threads, and the Linux aio-max-nr setting can

exceed system limits. Ideally, increase the `aio-max-nr` setting; as a workaround, you might reduce the settings for one or both of the MySQL configuration options.

Also take into consideration the value of `sync_binlog`, which controls synchronization of the binary log to disk.

For general I/O tuning advice, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

15.14 InnoDB INFORMATION_SCHEMA Tables

This section provides information and usage examples for `InnoDB INFORMATION_SCHEMA` tables.

`InnoDB INFORMATION_SCHEMA` tables provide metadata, status information, and statistics about various aspects of the `InnoDB` storage engine. You can view a list of `InnoDB INFORMATION_SCHEMA` tables by issuing a `SHOW TABLES` statement on the `INFORMATION_SCHEMA` database:

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA LIKE 'INNODB%';
```

For table definitions, see Section 24.31, “InnoDB INFORMATION_SCHEMA Tables”. For general information regarding the `MySQL INFORMATION_SCHEMA` database, see Chapter 24, *INFORMATION_SCHEMA Tables*.

15.14.1 InnoDB INFORMATION_SCHEMA Tables about Compression

There are two pairs of `InnoDB INFORMATION_SCHEMA` tables about compression that can provide insight into how well compression is working overall:

- `INNODB_CMP` and `INNODB_CMP_RESET` contain information about the number of compression operations and the amount of time spent performing compression.
- `INNODB_CMPMEM` and `INNODB_CMP_RESET` contain information about the way memory is allocated for compression.

15.14.1.1 INNODB_CMP and INNODB_CMP_RESET

The `INNODB_CMP` and `INNODB_CMP_RESET` tables contain status information about operations related to compressed tables, which are described in Section 15.9, “InnoDB Table and Page Compression”. The `PAGE_SIZE` column reports the compressed page size.

These two tables have identical contents, but reading from `INNODB_CMP_RESET` resets the statistics on compression and uncompression operations. For example, if you archive the output of `INNODB_CMP_RESET` every 60 minutes, you see the statistics for each hourly period. If you monitor the output of `INNODB_CMP` (making sure never to read `INNODB_CMP_RESET`), you see the cumulated statistics since InnoDB was started.

For the table definition, see Section 24.31.5, “The INFORMATION_SCHEMA INNODB_CMP and INNODB_CMP_RESET Tables”.

15.14.1.2 INNODB_CMPMEM and INNODB_CMPMEM_RESET

The `INNODB_CMPMEM` and `INNODB_CMPMEM_RESET` tables contain status information about compressed pages that reside in the buffer pool. Please consult Section 15.9, “InnoDB Table and Page Compression” for further information on compressed tables and the use of the buffer pool. The `INNODB_CMP` and `INNODB_CMP_RESET` tables should provide more useful statistics on compression.

Internal Details

`InnoDB` uses a buddy allocator system to manage memory allocated to pages of various sizes, from 1KB to 16KB. Each row of the two tables described here corresponds to a single page size.

The INNODB_CMPMEM and INNODB_CMPMEM_RESET tables have identical contents, but reading from INNODB_CMPMEM_RESET resets the statistics on relocation operations. For example, if every 60 minutes you archived the output of INNODB_CMPMEM_RESET, it would show the hourly statistics. If you never read INNODB_CMPMEM_RESET and monitored the output of INNODB_CMPMEM instead, it would show the cumulated statistics since InnoDB was started.

For the table definition, see Section 24.31.6, "The INFORMATION_SCHEMA INNODB_CMPMEM and INNODB_CMPMEM_RESET Tables".

15.14.1.3 Using the Compression Information Schema Tables

Example 15.1 Using the Compression Information Schema Tables

The following is sample output from a database that contains compressed tables (see Section 15.9, "InnoDB Table and Page Compression", INNODB_CMP, INNODB_CMP_PER_INDEX, and INNODB_CMPMEM).

The following table shows the contents of INFORMATION_SCHEMA.INNODB_CMP under a light workload. The only compressed page size that the buffer pool contains is 8K. Compressing or uncompressing pages has consumed less than a second since the time the statistics were reset, because the columns COMPRESS_TIME and UNCOMPRESS_TIME are zero.

page size	compress ops	compress ops ok	compress time	uncompress ops	uncompress time
1024	0	0	0	0	0
2048	0	0	0	0	0
4096	0	0	0	0	0
8192	1048	921	0	61	0
16384	0	0	0	0	0

According to INNODB_CMPMEM, there are 6169 compressed 8KB pages in the buffer pool. The only other allocated block size is 64 bytes. The smallest PAGE_SIZE in INNODB_CMPMEM is used for block descriptors of those compressed pages for which no uncompressed page exists in the buffer pool. We see that there are 5910 such pages. Indirectly, we see that 259 (6169-5910) compressed pages also exist in the buffer pool in uncompressed form.

The following table shows the contents of INFORMATION_SCHEMA.INNODB_CMPMEM under a light workload. Some memory is unusable due to fragmentation of the memory allocator for compressed pages: SUM(PAGE_SIZE*PAGES_FREE)=6784. This is because small memory allocation requests are fulfilled by splitting bigger blocks, starting from the 16K blocks that are allocated from the main buffer pool, using the buddy allocation system. The fragmentation is this low because some allocated blocks have been relocated (copied) to form bigger adjacent free blocks. This copying of SUM(PAGE_SIZE*RELOCATION_OPS) bytes has consumed less than a second (SUM(RELOCATION_TIME)=0).

page size	pages used	pages free	relocation ops	relocation time
64	5910	0	2436	0
128	0	1	0	0
256	0	0	0	0
512	0	1	0	0
1024	0	0	0	0
2048	0	1	0	0
4096	0	1	0	0
8192	6169	0	5	0
16384	0	0	0	0

15.14.2 InnoDB INFORMATION_SCHEMA Transaction and Locking Information

Note

This section describes locking information as exposed by the Performance Schema `data_locks` and `data_lock_waits` tables, which supersede the `INFORMATION_SCHEMA INNODB_LOCKS` and `INNODB_LOCK_WAITS` tables in MySQL 8.0. For similar discussion written in terms of the older `INFORMATION_SCHEMA` tables, see InnoDB INFORMATION_SCHEMA Transaction and Locking Information in MySQL 5.7 Reference Manual.

One `INFORMATION_SCHEMA` table and two Performance Schema tables enable you to monitor `InnoDB` transactions and diagnose potential locking problems:

- `INNODB_TRX`: This `INFORMATION_SCHEMA` table contains information about every transaction currently executing inside `InnoDB`, including the transaction state (for example, whether it is running or waiting for a lock), when the transaction started, and the particular SQL statement the transaction is executing.
- `data_locks`: This Performance Schema table contains a row for each hold lock and each lock request that is blocked waiting for a held lock to be released:
  - There is one row for each held lock, whatever the state of the transaction that holds the lock (`INNODB_TRX.TRX_STATE` is `RUNNING`, `LOCK WAIT`, `ROLLING BACK` or `COMMITTING`).
  - Each transaction in InnoDB that is waiting for another transaction to release a lock (`INNODB_TRX.TRX_STATE` is `LOCK WAIT`) is blocked by exactly one blocking lock request. That blocking lock request is for a row or table lock held by another transaction in an incompatible mode. A lock request always has a mode that is incompatible with the mode of the held lock that blocks the request (read vs. write, shared vs. exclusive).

    The blocked transaction cannot proceed until the other transaction commits or rolls back, thereby releasing the requested lock. For every blocked transaction, `data_locks` contains one row that describes each lock the transaction has requested, and for which it is waiting.
- `data_lock_waits`: This Performance Schema table indicates which transactions are waiting for a given lock, or for which lock a given transaction is waiting. This table contains one or more rows for each blocked transaction, indicating the lock it has requested and any locks that are blocking that request. The `REQUESTING_ENGINE_LOCK_ID` value refers to the lock requested by a transaction, and the `BLOCKING_ENGINE_LOCK_ID` value refers to the lock (held by another transaction) that prevents the first transaction from proceeding. For any given blocked transaction, all rows in `data_lock_waits` have the same value for `REQUESTING_ENGINE_LOCK_ID` and different values for `BLOCKING_ENGINE_LOCK_ID`.

For more information about the preceding tables, see Section 24.31.28, “The INFORMATION_SCHEMA INNODB_TRX Table”, Section 25.11.12.1, “The data_locks Table”, and Section 25.11.12.2, “The data_lock_waits Table”.

15.14.2.1 Using InnoDB Transaction and Locking Information

Note

This section describes locking information as exposed by the Performance Schema `data_locks` and `data_lock_waits` tables, which supersede the `INFORMATION_SCHEMA INNODB_LOCKS` and `INNODB_LOCK_WAITS` tables in MySQL 8.0. For similar discussion written in terms of the older `INFORMATION_SCHEMA` tables, see Using InnoDB Transaction and Locking Information in MySQL 5.7 Reference Manual.

Identifying Blocking Transactions

It is sometimes helpful to identify which transaction blocks another. The tables that contain information about `InnoDB` transactions and data locks enable you to determine which transaction is waiting for another, and which resource is being requested. (For descriptions of these tables, see Section 15.14.2, "InnoDB INFORMATION_SCHEMA Transaction and Locking Information".)

Suppose that three sessions are running concurrently. Each session corresponds to a MySQL thread, and executes one transaction after another. Consider the state of the system when these sessions have issued the following statements, but none has yet committed its transaction:

- Session A:

```
BEGIN;
SELECT a FROM t FOR UPDATE;
SELECT SLEEP(100);
```

- Session B:

```
SELECT b FROM t FOR UPDATE;
```

- Session C:

```
SELECT c FROM t FOR UPDATE;
```

In this scenario, use the following query to see which transactions are waiting and which transactions are blocking them:

```
SELECT
  r.trx_id waiting_trx_id,
  r.trx_mysql_thread_id waiting_thread,
  r.trx_query waiting_query,
  b.trx_id blocking_trx_id,
  b.trx_mysql_thread_id blocking_thread,
  b.trx_query blocking_query
FROM       performance_schema.data_lock_waits w
INNER JOIN information_schema.innodb_trx b
  ON b.trx_id = w.blocking_engine_transaction_id
INNER JOIN information_schema.innodb_trx r
  ON r.trx_id = w.requesting_engine_transaction_id;
```

Or, more simply, use the `sys` schema `innodb_lock_waits` view:

```
SELECT
  waiting_trx_id,
  waiting_pid,
  waiting_query,
  blocking_trx_id,
  blocking_pid,
  blocking_query
FROM sys.innodb_lock_waits;
```

waiting trx id	waiting thread	waiting query	blocking trx id	blocking thread	blocking query
A4	6	SELECT b FROM t FOR UPDATE	A3	5	SELECT SLEEP(100)
A5	7	SELECT c FROM t FOR UPDATE	A3	5	SELECT SLEEP(100)
A5	7	SELECT c FROM t FOR UPDATE	A4	6	SELECT b FROM t FOR UPDATE

In the preceding table, you can identify sessions by the “waiting query” or “blocking query” columns. As you can see:

- Session B (trx id A4, thread 6) and Session C (trx id A5, thread 7) are both waiting for Session A (trx id A3, thread 5).
- Session C is waiting for Session B as well as Session A.

You can see the underlying data in the INFORMATION_SCHEMA INNODB_TRX table and Performance Schema data_locks and data_lock_waits tables.

The following table shows some sample contents of the INNODB_TRX table.

trx id	trx state	trx started	trx requested lock id	trx wait started	trx weight	trx mysql thread id	trx query
A3	RUNNING	2008-01-15 16:44:54	NULL	NULL	2	5	SELECT SLEEP(100)
A4	LOCK WAIT	2008-01-15 16:45:09	A4:1:3:2	2008-01-15 16:45:09	2	6	SELECT b FROM t FOR UPDATE
A5	LOCK WAIT	2008-01-15 16:45:14	A5:1:3:2	2008-01-15 16:45:14	2	7	SELECT c FROM t FOR UPDATE

The following table shows some sample contents of the data_locks table.

lock id	lock trx id	lock mode	lock type	lock schema	lock table	lock index	lock data
A3:1:3:2	A3	X	RECORD	test	t	PRIMARY	0x0200
A4:1:3:2	A4	X	RECORD	test	t	PRIMARY	0x0200
A5:1:3:2	A5	X	RECORD	test	t	PRIMARY	0x0200

The following table shows some sample contents of the data_lock_waits table.

requesting trx id	requested lock id	blocking trx id	blocking lock id
A4	A4:1:3:2	A3	A3:1:3:2
A5	A5:1:3:2	A3	A3:1:3:2
A5	A5:1:3:2	A4	A4:1:3:2

Correlating InnoDB Transactions with MySQL Sessions

Sometimes it is useful to correlate internal InnoDB locking information with the session-level information maintained by MySQL. For example, you might like to know, for a given InnoDB transaction ID, the corresponding MySQL session ID and name of the session that may be holding a lock, and thus blocking other transactions.

The following output from the INFORMATION_SCHEMA INNODB_TRX table and Performance Schema data_locks and data_lock_waits tables is taken from a somewhat loaded system. As can be seen, there are several transactions running.

The following data_locks and data_lock_waits tables show that:

- Transaction 77F (executing an INSERT) is waiting for transactions 77E, 77D, and 77B to commit.
- Transaction 77E (executing an INSERT) is waiting for transactions 77D and 77B to commit.
- Transaction 77D (executing an INSERT) is waiting for transaction 77B to commit.
- Transaction 77B (executing an INSERT) is waiting for transaction 77A to commit.

- Transaction 77A is running, currently executing SELECT.
- Transaction E56 (executing an INSERT) is waiting for transaction E55 to commit.
- Transaction E55 (executing an INSERT) is waiting for transaction 19C to commit.
- Transaction 19C is running, currently executing an INSERT.

Note

There may be inconsistencies between queries shown in the INFORMATION_SCHEMA PROCESSLIST and INNODB_TRX tables. For an explanation, see Section 15.14.2.3, "Persistence and Consistency of InnoDB Transaction and Locking Information".

The following table shows the contents of the PROCESSLIST table for a system running a heavy workload.

ID	USER	HOST	DB	COMMAND	TIME	STATE	INFO
384	root	localhost	test	Query	10	update	INSERT INTO t2 VALUES …
257	root	localhost	test	Query	3	update	INSERT INTO t2 VALUES …
130	root	localhost	test	Query	0	update	INSERT INTO t2 VALUES …
61	root	localhost	test	Query	1	update	INSERT INTO t2 VALUES …
8	root	localhost	test	Query	1	update	INSERT INTO t2 VALUES …
4	root	localhost	test	Query	0	preparing	SELECT * FROM PROCESSLIST
2	root	localhost	test	Sleep	566		NULL

The following table shows the contents of the INNODB_TRX table for a system running a heavy workload.

trx id	trx state	trx started	trx requested lock id	trx wait started	trx weight	trx mysql thread id	trx query
77F	LOCK WAIT	2008-01-15 13:10:16	77F	2008-01-15 13:10:16	1	876	INSERT INTO t09 (D, B, C) VALUES …
77E	LOCK WAIT	2008-01-15 13:10:16	77E	2008-01-15 13:10:16	1	875	INSERT INTO t09 (D, B, C) VALUES …
77D	LOCK WAIT	2008-01-15 13:10:16	77D	2008-01-15 13:10:16	1	874	INSERT INTO t09 (D, B, C) VALUES …
77B	LOCK WAIT	2008-01-15 13:10:16	77B:733:12:1	2008-01-15 13:10:16	4	873	INSERT INTO t09 (D, B, C) VALUES …
77A	RUN-NING	2008-01-15 13:10:16	NULL	NULL	4	872	SELECT b, c FROM t09 WHERE …

trx id	trx state	trx started	trx requested lock id	trx wait started	trx weight	trx mysql thread id	trx query
E56	LOCK WAIT	2008-01-15 13:10:06	E56:743:6:2	2008-01-15 13:10:06	5	384	INSERT INTO t2 VALUES …
E55	LOCK WAIT	2008-01-15 13:10:06	E55:743:38:2	2008-01-15 13:10:13	965	257	INSERT INTO t2 VALUES …
19C	RUN-NING	2008-01-15 13:09:10	NULL	NULL	2900	130	INSERT INTO t2 VALUES …
E15	RUN-NING	2008-01-15 13:08:59	NULL	NULL	5395	61	INSERT INTO t2 VALUES …
51D	RUN-NING	2008-01-15 13:08:47	NULL	NULL	9807	8	INSERT INTO t2 VALUES …

The following table shows the contents of the `data_lock_waits` table for a system running a heavy workload.

requesting trx id	requested lock id	blocking trx id	blocking lock id
77F	77F:806	77E	77E:806
77F	77F:806	77D	77D:806
77F	77F:806	77B	77B:806
77E	77E:806	77D	77D:806
77E	77E:806	77B	77B:806
77D	77D:806	77B	77B:806
77B	77B:733:12:1	77A	77A:733:12:1
E56	E56:743:6:2	E55	E55:743:6:2
E55	E55:743:38:2	19C	19C:743:38:2

The following table shows the contents of the `data_locks` table for a system running a heavy workload.

lock id	lock trx id	lock mode	lock type	lock schema	lock table	lock index	lock data
77F:806	77F	AUTO_INC	TABLE	test	t09	NULL	NULL
77E:806	77E	AUTO_INC	TABLE	test	t09	NULL	NULL
77D:806	77D	AUTO_INC	TABLE	test	t09	NULL	NULL
77B:806	77B	AUTO_INC	TABLE	test	t09	NULL	NULL
77B:733:12:1	77B	X	RECORD	test	t09	PRIMARY	supremum pseudo-record
77A:733:12:1	77A	X	RECORD	test	t09	PRIMARY	supremum pseudo-record
E56:743:6:2	E56	S	RECORD	test	t2	PRIMARY	0, 0
E55:743:6:2	E55	X	RECORD	test	t2	PRIMARY	0, 0
E55:743:38:2	E55	S	RECORD	test	t2	PRIMARY	1922, 1922
19C:743:38:2	19C	X	RECORD	test	t2	PRIMARY	1922, 1922

15.14.2.2 InnoDB Lock and Lock-Wait Information

Note

This section describes locking information as exposed by the Performance Schema `data_locks` and `data_lock_waits` tables, which supersede the `INFORMATION_SCHEMA` `INNODB_LOCKS` and `INNODB_LOCK_WAITS` tables in MySQL 8.0. For similar discussion written in terms of the older `INFORMATION_SCHEMA` tables, see InnoDB Lock and Lock-Wait Information in MySQL 5.7 Reference Manual.

When a transaction updates a row in a table, or locks it with `SELECT FOR UPDATE`, `InnoDB` establishes a list or queue of locks on that row. Similarly, `InnoDB` maintains a list of locks on a table for table-level locks. If a second transaction wants to update a row or lock a table already locked by a prior transaction in an incompatible mode, `InnoDB` adds a lock request for the row to the corresponding queue. For a lock to be acquired by a transaction, all incompatible lock requests previously entered into the lock queue for that row or table must be removed (which occurs when the transactions holding or requesting those locks either commit or roll back).

A transaction may have any number of lock requests for different rows or tables. At any given time, a transaction may request a lock that is held by another transaction, in which case it is blocked by that other transaction. The requesting transaction must wait for the transaction that holds the blocking lock to commit or roll back. If a transaction is not waiting for a lock, it is in a `RUNNING` state. If a transaction is waiting for a lock, it is in a `LOCK WAIT` state. (The `INFORMATION_SCHEMA` `INNODB_TRX` table indicates transaction state values.)

The Performance Schema `data_locks` table holds one or more rows for each `LOCK WAIT` transaction, indicating any lock requests that prevent its progress. This table also contains one row describing each lock in a queue of locks pending for a given row or table. The Performance Schema `data_lock_waits` table shows which locks already held by a transaction are blocking locks requested by other transactions.

15.14.2.3 Persistence and Consistency of InnoDB Transaction and Locking Information

Note

This section describes locking information as exposed by the Performance Schema `data_locks` and `data_lock_waits` tables, which supersede the `INFORMATION_SCHEMA` `INNODB_LOCKS` and `INNODB_LOCK_WAITS` tables in MySQL 8.0. For similar discussion written in terms of the older `INFORMATION_SCHEMA` tables, see Persistence and Consistency of InnoDB Transaction and Locking Information in MySQL 5.7 Reference Manual.

The data exposed by the transaction and locking tables (`INFORMATION_SCHEMA` `INNODB_TRX` table, Performance Schema `data_locks` and `data_lock_waits` tables) represents a glimpse into fast-changing data. This is not like user tables, where the data changes only when application-initiated updates occur. The underlying data is internal system-managed data, and can change very quickly:

- Data might not be consistent between the `INNODB_TRX`, `data_locks`, and `data_lock_waits` tables.

  The `data_locks` and `data_lock_waits` tables expose live data from the `InnoDB` storage engine, to provide lock inormation about the transactions in the `INNODB_TRX` table. Data retrieved from the lock tables exists when the `SELECT` is executed, but might be gone or changed by the time the query result is consumed by the client.

  Joining `data_locks` with `data_lock_waits` can show rows in `data_lock_waits` that identify a parent row in `data_locks` that no longer exists or does not exist yet.

- Data in the transaction and locking tables might not be consistent with data in the `INFORMATION_SCHEMA` `PROCESSLIST` table or Performance Schema `threads` table.

For example, you should be careful when comparing data in the `InnoDB` transaction and locking tables with data in the `PROCESSLIST` table. Even if you issue a single `SELECT` (joining `INNODB_TRX` and `PROCESSLIST`, for example), the content of those tables is generally not consistent. It is possible for `INNODB_TRX` to reference rows that are not present in `PROCESSLIST` or for the currently executing SQL query of a transaction shown in `INNODB_TRX.TRX_QUERY` to differ from the one in `PROCESSLIST.INFO`.

15.14.3 InnoDB INFORMATION_SCHEMA System Tables

You can extract metadata about schema objects managed by `InnoDB` using `InnoDB INFORMATION_SCHEMA` system tables. This information comes from the data dictionary. Traditionally, you would get this type of information using the techniques from Section 15.16, “InnoDB Monitors”, setting up `InnoDB` monitors and parsing the output from the `SHOW ENGINE INNODB STATUS` statement. The `InnoDB INFORMATION_SCHEMA` table interface allows you to query this data using SQL.

With the exception of `INNODB_SYS_TABLESTATS`, for which there is no corresponding internal system table, `InnoDB INFORMATION_SCHEMA` system tables are populated with data read directly from internal `InnoDB` system tables rather than from metadata that is cached in memory.

`InnoDB INFORMATION_SCHEMA` system tables include the tables listed below. `INNODB_SYS_DATAFILES` and `INNODB_SYS_TABLESPACES` were added in with the introduction of support for the `DATA DIRECTORY='directory'` clause of the `CREATE TABLE` statement, which allows `InnoDB` file-per-table tablespaces (`.ibd` files) to be created in a location outside the MySQL data directory.

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA LIKE 'INNODB_SYS%';
+-----------------------------------------+
| Tables_in_information_schema (INNODB_SYS%) |
+-----------------------------------------+
| INNODB_SYS_DATAFILES                    |
| INNODB_SYS_TABLESTATS                   |
| INNODB_SYS_FOREIGN                      |
| INNODB_SYS_COLUMNS                      |
| INNODB_SYS_INDEXES                      |
| INNODB_SYS_FIELDS                       |
| INNODB_SYS_TABLESPACES                  |
| INNODB_SYS_FOREIGN_COLS                 |
| INNODB_SYS_TABLES                       |
+-----------------------------------------+
```

The table names are indicative of the type of data provided:

- `INNODB_SYS_TABLES` provides metadata about `InnoDB` tables.
- `INNODB_SYS_COLUMNS` provides metadata about `InnoDB` table columns.
- `INNODB_SYS_INDEXES` provides metadata about `InnoDB` indexes.
- `INNODB_SYS_FIELDS` provides metadata about the key columns (fields) of `InnoDB` indexes.
- `INNODB_SYS_TABLESTATS` provides a view of low-level status information about `InnoDB` tables that is derived from in-memory data structures.
- `INNODB_SYS_DATAFILES` provides data file path information for `InnoDB` file-per-table and general tablespaces.
- `INNODB_SYS_TABLESPACES` provides metadata about `InnoDB` file-per-table and general tablespaces.
- `INNODB_SYS_FOREIGN` provides metadata about foreign keys defined on `InnoDB` tables.
- `INNODB_SYS_FOREIGN_COLS` provides metadata about the columns of foreign keys that are defined on `InnoDB` tables.

InnoDB INFORMATION_SCHEMA system tables can be joined together through fields such as TABLE_ID, INDEX_ID, and SPACE, allowing you to easily retrieve all available data for an object you want to study or monitor.

Refer to the InnoDB INFORMATION_SCHEMA documentation for information about the columns of each table.

Example 15.2 InnoDB INFORMATION_SCHEMA System Tables

This example uses a simple table (t1) with a single index (i1) to demonstrate the type of metadata found in the InnoDB INFORMATION_SCHEMA system tables.

1. Create a test database and table t1:

```
mysql> CREATE DATABASE test;

mysql> USE test;

mysql> CREATE TABLE t1 (
col1 INT,
col2 CHAR(10),
col3 VARCHAR(10))
ENGINE = InnoDB;

mysql> CREATE INDEX i1 ON t1(col1);
```

2. After creating the table t1, query INNODB_SYS_TABLES to locate the metadata for test/t1:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME='test/t1' \G

*************************** 1. row ***************************
     TABLE_ID: 71
         NAME: test/t1
         FLAG: 1
       N_COLS: 6
        SPACE: 57
   ROW_FORMAT: Compact
ZIP_PAGE_SIZE: 0
...
```

   Table t1 has a TABLE_ID of 71. The FLAG field provides bit level information about table format and storage characteristics. There are six columns, three of which are hidden columns created by InnoDB (DB_ROW_ID, DB_TRX_ID, and DB_ROLL_PTR). The ID of the table's SPACE is 57 (a value of 0 would indicate that the table resides in the system tablespace). The ROW_FORMAT is Compact. ZIP_PAGE_SIZE only applies to tables with a Compressed row format.

3. Using the TABLE_ID information from INNODB_SYS_TABLES, query the INNODB_SYS_COLUMNS table for information about the table's columns.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_COLUMNS where TABLE_ID = 71 \G
*************************** 1. row ***************************
TABLE_ID: 71
    NAME: col1
     POS: 0
   MTYPE: 6
  PRTYPE: 1027
     LEN: 4
*************************** 2. row ***************************
TABLE_ID: 71
    NAME: col2
     POS: 1
   MTYPE: 2
  PRTYPE: 524542
     LEN: 10
*************************** 3. row ***************************
TABLE_ID: 71
    NAME: col3
     POS: 2
   MTYPE: 1
```

```
  PRTYPE: 524303
     LEN: 10
```

In addition to the TABLE_ID and column NAME, INNODB_SYS_COLUMNS provides the ordinal position (POS) of each column (starting from 0 and incrementing sequentially), the column MTYPE or “main type” (6 = INT, 2 = CHAR, 1 = VARCHAR), the PRTYPE or “precise type” (a binary value with bits that represent the MySQL data type, character set code, and nullability), and the column length (LEN).

4. Using the TABLE_ID information from INNODB_SYS_TABLES once again, query INNODB_SYS_INDEXES for information about the indexes associated with table t1.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_INDEXES WHERE TABLE_ID = 71 \G
*************************** 1. row ***************************
       INDEX_ID: 111
           NAME: GEN_CLUST_INDEX
       TABLE_ID: 71
           TYPE: 1
       N_FIELDS: 0
        PAGE_NO: 3
          SPACE: 57
MERGE_THRESHOLD: 50
*************************** 2. row ***************************
       INDEX_ID: 112
           NAME: i1
       TABLE_ID: 71
           TYPE: 0
       N_FIELDS: 1
        PAGE_NO: 4
          SPACE: 57
MERGE_THRESHOLD: 50
```

INNODB_SYS_INDEXES returns data for two indexes. The first index is GEN_CLUST_INDEX, which is a clustered index created by InnoDB if the table does not have a user-defined clustered index. The second index (i1) is the user-defined secondary index.

The INDEX_ID is an identifier for the index that is unique across all databases in an instance. The TABLE_ID identifies the table that the index is associated with. The index TYPE value indicates the type of index (1 = Clustered Index, 0 = Secondary index). The N_FILEDS value is the number of fields that comprise the index. PAGE_NO is the root page number of the index B-tree, and SPACE is the ID of the tablespace where the index resides. A non-zero value indicates that the index does not reside in the system tablespace. MERGE_THRESHOLD defines a percentage threshold value for the amount of data in an index page. If the amount of data in an index page falls below the this value (the default is 50%) when a row is deleted or when a row is shortened by an update operation, InnoDB attempts to merge the index page with a neighboring index page.

5. Using the INDEX_ID information from INNODB_SYS_INDEXES, query INNODB_SYS_FIELDS for information about the fields of index i1.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FIELDS where INDEX_ID = 112 \G
*************************** 1. row ***************************
INDEX_ID: 112
    NAME: col1
     POS: 0
```

INNODB_SYS_FIELDS provides the NAME of the indexed field and its ordinal position within the index. If the index (i1) had been defined on multiple fields, INNODB_SYS_FIELDS would provide metadata for each of the indexed fields.

6. Using the SPACE information from INNODB_SYS_TABLES, query INNODB_SYS_TABLESPACES table for information about the table's tablespace.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES WHERE SPACE = 57 \G
*************************** 1. row ***************************
       SPACE: 57
        NAME: test/t1
```

```
         FLAG: 0
   ROW_FORMAT: Compact or Redundant
    PAGE_SIZE: 16384
ZIP_PAGE_SIZE: 0
```

In addition to the SPACE ID of the tablespace and the NAME of the associated table, INNODB_SYS_TABLESPACES provides tablespace FLAG data, which is bit level information about tablespace format and storage characteristics. Also provided are tablespace ROW_FORMAT, PAGE_SIZE, and ZIP_PAGE_SIZE data (ZIP_PAGE_SIZE is applicable to tablespaces with a Compressed row format).

7. Using the SPACE information from INNODB_SYS_TABLES once again, query INNODB_SYS_DATAFILES for the location of the tablespace data file.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_DATAFILES WHERE SPACE = 57 \G
*************************** 1. row ***************************
SPACE: 57
 PATH: ./test/t1.ibd
```

The datafile is located in the test directory under MySQL's data directory. If a file-per-table tablespace were created in a location outside the MySQL data directory using the DATA DIRECTORY clause of the CREATE TABLE statement, the tablespace PATH would be a fully qualified directory path.

8. As a final step, insert a row into table t1 (TABLE_ID = 71) and view the data in the INNODB_SYS_TABLESTATS table. The data in this table is used by the MySQL optimizer to calculate which index to use when querying an InnoDB table. This information is derived from in-memory data structures. There is no corresponding internal InnoDB system table.

```
mysql> INSERT INTO t1 VALUES(5, 'abc', 'def');
Query OK, 1 row affected (0.06 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLESTATS where TABLE_ID = 71 \G
*************************** 1. row ***************************
         TABLE_ID: 71
             NAME: test/t1
STATS_INITIALIZED: Initialized
         NUM_ROWS: 1
 CLUST_INDEX_SIZE: 1
 OTHER_INDEX_SIZE: 0
 MODIFIED_COUNTER: 1
          AUTOINC: 0
        REF_COUNT: 1
```

The STATS_INITIALIZED field indicates whether or not statistics have been collected for the table. NUM_ROWS is the current estimated number of rows in the table. The CLUST_INDEX_SIZE and OTHER_INDEX_SIZE fields report the number of pages on disk that store clustered and secondary indexes for the table, respectively. The MODIFIED_COUNTER value shows the number of rows modified by DML operations and cascade operations from foreign keys. The AUTOINC value is the next number to be issued for any autoincrement-based operation. There are no autoincrement columns defined on table t1, so the value is 0. The REF_COUNT value is a counter. When the counter reaches 0, it signifies that the table metadata can be evicted from the table cache.

Example 15.3 Foreign Key INFORMATION_SCHEMA System Tables

The INNODB_SYS_FOREIGN and INNODB_SYS_FOREIGN_COLS tables provide data about foreign key relationships. This example uses a parent table and child table with a foreign key relationship to demonstrate the data found in the INNODB_SYS_FOREIGN and INNODB_SYS_FOREIGN_COLS tables.

1. Create the test database with parent and child tables:

```
mysql> CREATE DATABASE test;

mysql> USE test;

mysql> CREATE TABLE parent (id INT NOT NULL,
```

```
    -> PRIMARY KEY (id)) ENGINE=INNODB;

mysql> CREATE TABLE child (id INT, parent_id INT,
    -> INDEX par_ind (parent_id),
    -> CONSTRAINT fk1
    -> FOREIGN KEY (parent_id) REFERENCES parent(id)
    -> ON DELETE CASCADE) ENGINE=INNODB;
```

2. After the parent and child tables are created, query `INNODB_SYS_FOREIGN` and locate the foreign key data for the `test/child` and `test/parent` foreign key relationship:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FOREIGN \G
*************************** 1. row ***************************
      ID: test/fk1
FOR_NAME: test/child
REF_NAME: test/parent
  N_COLS: 1
    TYPE: 1
```

   Metadata includes the foreign key `ID` (`fk1`), which is named for the `CONSTRAINT` that was defined on the child table. The `FOR_NAME` is the name of the child table where the foreign key is defined. `REF_NAME` is the name of the parent table (the “referenced” table). `N_COLS` is the number of columns in the foreign key index. `TYPE` is a numerical value representing bit flags that provide additional information about the foreign key column. In this case, the `TYPE` value is 1, which indicates that the `ON DELETE CASCADE` option was specified for the foreign key. See the `INNODB_SYS_FOREIGN` table definition for more information about `TYPE` values.

3. Using the foreign key `ID`, query `INNODB_SYS_FOREIGN_COLS` to view data about the columns of the foreign key.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FOREIGN_COLS WHERE ID = 'test/fk1' \G
*************************** 1. row ***************************
         ID: test/fk1
FOR_COL_NAME: parent_id
REF_COL_NAME: id
         POS: 0
```

   `FOR_COL_NAME` is the name of the foreign key column in the child table, and `REF_COL_NAME` is the name of the referenced column in the parent table. The `POS` value is the ordinal position of the key field within the foreign key index, starting at zero.

Example 15.4 Joining InnoDB INFORMATION_SCHEMA System Tables

This example demonstrates joining three `InnoDB INFORMATION_SCHEMA` system tables (`INNODB_SYS_TABLES`, `INNODB_SYS_TABLESPACES`, and `INNODB_SYS_TABLESTATS`) to gather file format, row format, page size, and index size information about tables in the employees sample database.

The following table name aliases are used to shorten the query string:

- `INFORMATION_SCHEMA.INNODB_SYS_TABLES`: a
- `INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES`: b
- `INFORMATION_SCHEMA.INNODB_SYS_TABLESTATS`: c

An `IF()` control flow function is used to account for compressed tables. If a table is compressed, the index size is calculated using `ZIP_PAGE_SIZE` rather than `PAGE_SIZE`. `CLUST_INDEX_SIZE` and `OTHER_INDEX_SIZE`, which are reported in bytes, are divided by `1024*1024` to provide index sizes in megabytes (MBs). MB values are rounded to zero decimal spaces using the `ROUND()` function.

```
mysql> SELECT a.NAME, a.ROW_FORMAT,
 @page_size :=
  IF(a.ROW_FORMAT='Compressed',
   b.ZIP_PAGE_SIZE, b.PAGE_SIZE)
   AS page_size,
```

```
  ROUND((@page_size * c.CLUST_INDEX_SIZE)
   /(1024*1024)) AS pk_mb,
  ROUND((@page_size * c.OTHER_INDEX_SIZE)
   /(1024*1024)) AS secidx_mb
FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES a
INNER JOIN INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES b on a.NAME = b.NAME
INNER JOIN INFORMATION_SCHEMA.INNODB_SYS_TABLESTATS c on b.NAME = c.NAME
WHERE a.NAME LIKE 'employees/%'
ORDER BY a.NAME DESC;
+-----------------------+------------+-----------+-------+-----------+
| NAME                  | ROW_FORMAT | page_size | pk_mb | secidx_mb |
+-----------------------+------------+-----------+-------+-----------+
| employees/titles      | Dynamic    |     16384 |    20 |        11 |
| employees/salaries    | Dynamic    |     16384 |    93 |        34 |
| employees/employees   | Dynamic    |     16384 |    15 |         0 |
| employees/dept_manager | Dynamic    |     16384 |     0 |         0 |
| employees/dept_emp    | Dynamic    |     16384 |    12 |        10 |
| employees/departments | Dynamic    |     16384 |     0 |         0 |
+-----------------------+------------+-----------+-------+-----------+
```

15.14.4 InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables

The following tables store metadata for `FULLTEXT` indexes:

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA LIKE 'INNODB_FT%';
+-----------------------------------------+
| Tables_in_INFORMATION_SCHEMA (INNODB_FT%) |
+-----------------------------------------+
| INNODB_FT_CONFIG                        |
| INNODB_FT_BEING_DELETED                 |
| INNODB_FT_DELETED                       |
| INNODB_FT_DEFAULT_STOPWORD              |
| INNODB_FT_INDEX_TABLE                   |
| INNODB_FT_INDEX_CACHE                   |
+-----------------------------------------+
```

Table Overview

- `INNODB_FT_CONFIG`: Displays metadata about the `FULLTEXT` index and associated processing for an `InnoDB` table.
- `INNODB_FT_BEING_DELETED`: Provides a snapshot of the `INNODB_FT_DELETED` table that is only used during an `OPTIMIZE TABLE` maintenance operation. When `OPTIMIZE TABLE` is run, the `INNODB_FT_BEING_DELETED` table is emptied, and DOC_IDs are removed from the `INNODB_FT_DELETED` table. Because the contents of `INNODB_FT_BEING_DELETED` typically have a short lifetime, this table has limited utility for monitoring or debugging. For information about running `OPTIMIZE TABLE` on tables with `FULLTEXT` indexes, see Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.
- `INNODB_FT_DELETED`: Records rows that are deleted from the `FULLTEXT` index for an `InnoDB` table. To avoid expensive index reorganization during DML operations for an `InnoDB` `FULLTEXT` index, the information about newly deleted words is stored separately, filtered out of search results when you perform a text search, and removed from the main search index only when you run `OPTIMIZE TABLE`.
- `INNODB_FT_DEFAULT_STOPWORD`: Holds a list of stopwords that are used by default when creating a `FULLTEXT` index.

  For information about the `INNODB_FT_DEFAULT_STOPWORD` table, see Section 12.9.4, “Full-Text Stopwords”.
- `INNODB_FT_INDEX_TABLE`: Contains data about the inverted index used to process text searches against the `FULLTEXT` index.
- `INNODB_FT_INDEX_CACHE`: Contains token information about newly inserted rows in a `FULLTEXT` index. To avoid expensive index reorganization during DML operations, the information about newly

indexed words is stored separately, and combined with the main search index only when `OPTIMIZE TABLE` is run, when the server is shut down, or when the cache size exceeds a limit defined by `innodb_ft_cache_size` or `innodb_ft_total_cache_size`.

Note

With the exception of the `INNODB_FT_DEFAULT_STOPWORD` table, you must set the `innodb_ft_aux_table` configuration variable to the name of the table (*`database_name`*/*`table_name`*) that contains the `FULLTEXT` index. Otherwise, the `InnoDB` `FULLTEXT` index `INFORMATION_SCHEMA` tables appear empty.

Example 15.5 InnoDB FULLTEXT Index INFORMATION_SCHEMA Tables

This example uses a table with a `FULLTEXT` index to demonstrate the data contained in the `FULLTEXT` index `INFORMATION_SCHEMA` tables.

1. Create a table with a `FULLTEXT` index and insert some data:

```
mysql> CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT,
      FULLTEXT (title,body)
     ) ENGINE=InnoDB;

INSERT INTO articles (title,body) VALUES
    ('MySQL Tutorial','DBMS stands for DataBase ...'),
    ('How To Use MySQL Well','After you went through a ...'),
    ('Optimizing MySQL','In this tutorial we will show ...'),
    ('1001 MySQL Tricks','1. Never run mysqld as root. 2. ...'),
    ('MySQL vs. YourSQL','In the following database comparison ...'),
    ('MySQL Security','When configured properly, MySQL ...');
```

2. Set the `innodb_ft_aux_table` variable to the name of the table with the `FULLTEXT` index. If this variable is not set, the `InnoDB` `FULLTEXT` `INFORMATION_SCHEMA` tables appear empty, with the exception of the `INNODB_FT_DEFAULT_STOPWORD` table.

```
SET GLOBAL innodb_ft_aux_table = 'test/articles';
```

3. Query the `INNODB_FT_INDEX_CACHE` table, which shows information about newly inserted rows in a `FULLTEXT` index. To avoid expensive index reorganization during DML operations, data for newly inserted rows remains in the `FULLTEXT` index cache until `OPTIMIZE TABLE` is run (or until the server is shutdown or cache limits are exceeded).

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE LIMIT 5;
+------------+-------------+-------------+-----------+--------+----------+
| WORD       | FIRST_DOC_ID | LAST_DOC_ID | DOC_COUNT | DOC_ID | POSITION |
+------------+-------------+-------------+-----------+--------+----------+
| 1001       |            5 |           5 |         1 |      5 |        0 |
| after      |            3 |           3 |         1 |      3 |       22 |
| comparison |            6 |           6 |         1 |      6 |       44 |
| configured |            7 |           7 |         1 |      7 |       20 |
| database   |            2 |           6 |         2 |      2 |       31 |
+------------+-------------+-------------+-----------+--------+----------+
```

4. Enable `innodb_optimize_fulltext_only` and run `OPTIMIZE TABLE` on the table that contains the `FULLTEXT` index. This operation flushes the contents of the `FULLTEXT` index cache to the main `FULLTEXT` index. `innodb_optimize_fulltext_only` changes the way the `OPTIMIZE TABLE` statement operates on `InnoDB` tables, and is intended to be enabled temporarily, during maintenance operations on `InnoDB` tables with `FULLTEXT` indexes.

```
mysql> SET GLOBAL innodb_optimize_fulltext_only=ON;
Query OK, 0 rows affected (0.00 sec)

mysql> OPTIMIZE TABLE articles;
+--------------+----------+----------+----------+
```

```
| Table         | Op       | Msg_type | Msg_text |
+---------------+----------+----------+----------+
| test.articles | optimize | status   | OK       |
+---------------+----------+----------+----------+
```

5. Query the `INNODB_FT_INDEX_TABLE` table to view information about data in the main `FULLTEXT` index, including information about the data that was just flushed from the `FULLTEXT` index cache.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_TABLE LIMIT 5;
+------------+--------------+-------------+-----------+--------+----------+
| WORD       | FIRST_DOC_ID | LAST_DOC_ID | DOC_COUNT | DOC_ID | POSITION |
+------------+--------------+-------------+-----------+--------+----------+
| 1001       |            5 |           5 |         1 |      5 |        0 |
| after      |            3 |           3 |         1 |      3 |       22 |
| comparison |            6 |           6 |         1 |      6 |       44 |
| configured |            7 |           7 |         1 |      7 |       20 |
| database   |            2 |           6 |         2 |      2 |       31 |
+------------+--------------+-------------+-----------+--------+----------+
```

   The `INNODB_FT_INDEX_CACHE` table is now empty since the `OPTIMIZE TABLE` operation flushed the `FULLTEXT` index cache.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_INDEX_CACHE LIMIT 5;
Empty set (0.00 sec)
```

6. Delete some records from the `test/articles` table.

```
mysql> DELETE FROM test.articles WHERE id < 4;
Query OK, 3 rows affected (0.11 sec)
```

7. Query the `INNODB_FT_DELETED` table. This table records rows that are deleted from the `FULLTEXT` index. To avoid expensive index reorganization during DML operations, information about newly deleted records is stored separately, filtered out of search results when you do a text search, and removed from the main search index when you run `OPTIMIZE TABLE`.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_DELETED;
+--------+
| DOC_ID |
+--------+
|      2 |
|      3 |
|      4 |
+--------+
```

8. Run `OPTIMIZE TABLE` to remove the deleted records.

```
mysql> OPTIMIZE TABLE articles;
+---------------+----------+----------+----------+
| Table         | Op       | Msg_type | Msg_text |
+---------------+----------+----------+----------+
| test.articles | optimize | status   | OK       |
+---------------+----------+----------+----------+
```

   The `INNODB_FT_DELETED` table should now appear empty.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_DELETED;
Empty set (0.00 sec)
```

9. Query the `INNODB_FT_CONFIG` table. This table contains metadata about the `FULLTEXT` index and related processing:

   - `optimize_checkpoint_limit` is the number of seconds after which an `OPTIMIZE TABLE` run stops.
   - `synced_doc_id` is the next `DOC_ID` to be issued.
   - `stopword_table_name` is the *`database/table`* name for a user-defined stopword table. This field appears empty if there is no user-defined stopword table.

- `use_stopword` indicates whether or not a stopword table is used, which is defined when the `FULLTEXT` index is created.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_CONFIG;
+---------------------------+-------+
| KEY                       | VALUE |
+---------------------------+-------+
| optimize_checkpoint_limit | 180   |
| synced_doc_id             | 8     |
| stopword_table_name       |       |
| use_stopword              | 1     |
+---------------------------+-------+
```

15.14.5 InnoDB INFORMATION_SCHEMA Buffer Pool Tables

The `InnoDB` `INFORMATION_SCHEMA` buffer pool tables provide buffer pool status information and metadata about the pages within the `InnoDB` buffer pool.

The `InnoDB` `INFORMATION_SCHEMA` buffer pool tables include those listed below:

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA LIKE 'INNODB_BUFFER%';
+-----------------------------------------------+
| Tables_in_INFORMATION_SCHEMA (INNODB_BUFFER%) |
+-----------------------------------------------+
| INNODB_BUFFER_PAGE_LRU                        |
| INNODB_BUFFER_PAGE                            |
| INNODB_BUFFER_POOL_STATS                      |
+-----------------------------------------------+
```

Table Overview

- `INNODB_BUFFER_PAGE`: Holds information about each page in the `InnoDB` buffer pool.
- `INNODB_BUFFER_PAGE_LRU`: Holds information about the pages in the `InnoDB` buffer pool, in particular how they are ordered in the LRU list that determines which pages to evict from the buffer pool when it becomes full. The `INNODB_BUFFER_PAGE_LRU` table has the same columns as the `INNODB_BUFFER_PAGE` table, except that the `INNODB_BUFFER_PAGE_LRU` table has an `LRU_POSITION` column instead of a `BLOCK_ID` column.
- `INNODB_BUFFER_POOL_STATS`: Provides buffer pool status information. Much of the same information is provided by `SHOW ENGINE INNODB STATUS` output, or may be obtained using `InnoDB` buffer pool server status variables.

Warning

Querying the `INNODB_BUFFER_PAGE` table or `INNODB_BUFFER_PAGE_LRU` table can introduce significant performance overhead. Do not query these tables on a production system unless you are aware of the performance impact that your query may have, and have determined it to be acceptable. To avoid impacting performance, reproduce the issue you want to investigate on a test instance and run your queries on the test instance.

Example 15.6 Querying System Data in the INNODB_BUFFER_PAGE Table

This query provides an approximate count of pages that contain system data by excluding pages where the `TABLE_NAME` value is either `NULL` or includes a slash `/` or period `.` in the table name, which indicates a user-defined table.

```
SELECT COUNT(*) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NULL OR (INSTR(TABLE_NAME, '/') = 0 AND INSTR(TABLE_NAME, '.') = 0);
+----------+
| COUNT(*) |
+----------+
|     1516 |
+----------+
```

This query returns the approximate number of pages that contain system data, the total number of buffer pool pages, and an approximate percentage of pages that contain system data.

```
SELECT
(SELECT COUNT(*) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NULL OR (INSTR(TABLE_NAME, '/') = 0 AND INSTR(TABLE_NAME, '.') = 0)
) AS system_pages,
(
SELECT COUNT(*)
FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
) AS total_pages,
(
SELECT ROUND((system_pages/total_pages) * 100)
) AS system_page_percentage;
+--------------+-------------+-----------------------+
| system_pages | total_pages | system_page_percentage |
+--------------+-------------+-----------------------+
|          295 |        8192 |                     4 |
+--------------+-------------+-----------------------+
```

The type of system data in the buffer pool can be determined by querying the `PAGE_TYPE` value. For example, the following query returns eight distinct `PAGE_TYPE` values among the pages that contain system data:

```
mysql> SELECT DISTINCT PAGE_TYPE FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NULL OR (INSTR(TABLE_NAME, '/') = 0 AND INSTR(TABLE_NAME, '.') = 0);
+-------------------+
| PAGE_TYPE         |
+-------------------+
| SYSTEM            |
| IBUF_BITMAP       |
| UNKNOWN           |
| FILE_SPACE_HEADER |
| INODE             |
| UNDO_LOG          |
| ALLOCATED         |
+-------------------+
```

Example 15.7 Querying User Data in the INNODB_BUFFER_PAGE Table

This query provides an approximate count of pages containing user data by counting pages where the `TABLE_NAME` value is `NOT NULL` and `NOT LIKE '%INNODB_SYS_TABLES%'`.

```
mysql> SELECT COUNT(*) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NOT NULL AND TABLE_NAME NOT LIKE '%INNODB_SYS_TABLES%';
+----------+
| COUNT(*) |
+----------+
|     7897 |
+----------+
```

This query returns the approximate number of pages that contain user data, the total number of buffer pool pages, and an approximate percentage of pages that contain user data.

```
mysql> SELECT
(SELECT COUNT(*) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NOT NULL AND (INSTR(TABLE_NAME, '/') > 0 OR INSTR(TABLE_NAME, '.') > 0)
) AS user_pages,
(
SELECT COUNT(*)
FROM information_schema.INNODB_BUFFER_PAGE
) AS total_pages,
(
SELECT ROUND((user_pages/total_pages) * 100)
) AS user_page_percentage;
+------------+-------------+----------------------+
| user_pages | total_pages | user_page_percentage |
+------------+-------------+----------------------+
|       7897 |        8192 |                   96 |
+------------+-------------+----------------------+
```

This query identifies user-defined tables with pages in the buffer pool:

```
mysql> SELECT DISTINCT TABLE_NAME FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME IS NOT NULL AND (INSTR(TABLE_NAME, '/') > 0 OR INSTR(TABLE_NAME, '.') > 0)
AND TABLE_NAME NOT LIKE '`mysql`.`innodb_%';
+-------------------------+
| TABLE_NAME              |
+-------------------------+
| `employees`.`salaries`  |
| `employees`.`employees` |
+-------------------------+
```

Example 15.8 Querying Index Data in the INNODB_BUFFER_PAGE Table

For information about index pages, query the `INDEX_NAME` column using the name of the index. For example, the following query returns the number of pages and total data size of pages for the `emp_no` index that is defined on the `employees.salaries` table:

```
mysql> SELECT INDEX_NAME, COUNT(*) AS Pages,
ROUND(SUM(IF(COMPRESSED_SIZE = 0, @@global.innodb_page_size, COMPRESSED_SIZE))/1024/1024)
AS 'Total Data (MB)'
FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE INDEX_NAME='emp_no' AND TABLE_NAME = '`employees`.`salaries`';
+------------+-------+-----------------+
| INDEX_NAME | Pages | Total Data (MB) |
+------------+-------+-----------------+
| emp_no     |  1609 |              25 |
+------------+-------+-----------------+
```

This query returns the number of pages and total data size of pages for all indexes defined on the `employees.salaries` table:

```
mysql> SELECT INDEX_NAME, COUNT(*) AS Pages,
ROUND(SUM(IF(COMPRESSED_SIZE = 0, @@global.innodb_page_size, COMPRESSED_SIZE))/1024/1024)
AS 'Total Data (MB)'
FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE
WHERE TABLE_NAME = '`employees`.`salaries`'
GROUP BY INDEX_NAME;
+------------+-------+-----------------+
| INDEX_NAME | Pages | Total Data (MB) |
+------------+-------+-----------------+
| emp_no     |  1608 |              25 |
| PRIMARY    |  6086 |              95 |
+------------+-------+-----------------+
```

Example 15.9 Querying LRU_POSITION Data in the INNODB_BUFFER_PAGE_LRU Table

The `INNODB_BUFFER_PAGE_LRU` table holds information about the pages in the `InnoDB` buffer pool, in particular how they are ordered that determines which pages to evict from the buffer pool when it becomes full. The definition for this page is the same as for `INNODB_BUFFER_PAGE`, except this table has an `LRU_POSITION` column instead of a `BLOCK_ID` column.

This query counts the number of positions at a specific location in the LRU list occupied by pages of the `employees.employees` table.

```
mysql> SELECT COUNT(LRU_POSITION) FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE_LRU
WHERE TABLE_NAME='`employees`.`employees`' AND LRU_POSITION < 3072;
+---------------------+
| COUNT(LRU_POSITION) |
+---------------------+
|                 548 |
+---------------------+
```

Example 15.10 Querying the INNODB_BUFFER_POOL_STATS Table

The `INNODB_BUFFER_POOL_STATS` table provides information similar to `SHOW ENGINE INNODB STATUS` and `InnoDB` buffer pool status variables.

```
mysql> SELECT * FROM information_schema.INNODB_BUFFER_POOL_STATS \G
*************************** 1. row ***************************
```

```
                        POOL_ID: 0
                      POOL_SIZE: 8192
                   FREE_BUFFERS: 1
                 DATABASE_PAGES: 8173
             OLD_DATABASE_PAGES: 3014
       MODIFIED_DATABASE_PAGES: 0
            PENDING_DECOMPRESS: 0
                 PENDING_READS: 0
             PENDING_FLUSH_LRU: 0
            PENDING_FLUSH_LIST: 0
              PAGES_MADE_YOUNG: 15907
          PAGES_NOT_MADE_YOUNG: 3803101
         PAGES_MADE_YOUNG_RATE: 0
     PAGES_MADE_NOT_YOUNG_RATE: 0
             NUMBER_PAGES_READ: 3270
          NUMBER_PAGES_CREATED: 13176
          NUMBER_PAGES_WRITTEN: 15109
               PAGES_READ_RATE: 0
             PAGES_CREATE_RATE: 0
            PAGES_WRITTEN_RATE: 0
              NUMBER_PAGES_GET: 33069332
                      HIT_RATE: 0
  YOUNG_MAKE_PER_THOUSAND_GETS: 0
NOT_YOUNG_MAKE_PER_THOUSAND_GETS: 0
       NUMBER_PAGES_READ_AHEAD: 2713
     NUMBER_READ_AHEAD_EVICTED: 0
            READ_AHEAD_RATE: 0
      READ_AHEAD_EVICTED_RATE: 0
                 LRU_IO_TOTAL: 0
               LRU_IO_CURRENT: 0
             UNCOMPRESS_TOTAL: 0
          UNCOMPRESS_CURRENT: 0
```

For comparison, `SHOW ENGINE INNODB STATUS` output and `InnoDB` buffer pool status variable output is shown below, based on the same data set.

For more information about `SHOW ENGINE INNODB STATUS` output, see Section 15.16.3, “InnoDB Standard Monitor and Lock Monitor Output”.

```
mysql> SHOW ENGINE INNODB STATUS \G
...
----------------------
BUFFER POOL AND MEMORY
----------------------
Total large memory allocated 137428992
Dictionary memory allocated 579084
Buffer pool size   8192
Free buffers       1
Database pages     8173
Old database pages 3014
Modified db pages  0
Pending reads 0
Pending writes: LRU 0, flush list 0, single page 0
Pages made young 15907, not young 3803101
0.00 youngs/s, 0.00 non-youngs/s
Pages read 3270, created 13176, written 15109
0.00 reads/s, 0.00 creates/s, 0.00 writes/s
No buffer pool page gets since the last printout
Pages read ahead 0.00/s, evicted without access 0.00/s, Random read ahead 0.00/s
LRU len: 8173, unzip_LRU len: 0
I/O sum[0]:cur[0], unzip sum[0]:cur[0]
...
```

For status variable descriptions, see Section 5.1.7, “Server Status Variables”.

```
mysql> SHOW STATUS LIKE 'Innodb_buffer%';
+---------------------------------------+-------------+
| Variable_name                         | Value       |
+---------------------------------------+-------------+
| Innodb_buffer_pool_dump_status        | not started |
| Innodb_buffer_pool_load_status        | not started |
| Innodb_buffer_pool_resize_status      | not started |
```

```
| Innodb_buffer_pool_pages_data         | 8173        |
| Innodb_buffer_pool_bytes_data         | 133906432   |
| Innodb_buffer_pool_pages_dirty        | 0           |
| Innodb_buffer_pool_bytes_dirty        | 0           |
| Innodb_buffer_pool_pages_flushed      | 15109       |
| Innodb_buffer_pool_pages_free         | 1           |
| Innodb_buffer_pool_pages_misc         | 18          |
| Innodb_buffer_pool_pages_total        | 8192        |
| Innodb_buffer_pool_read_ahead_rnd     | 0           |
| Innodb_buffer_pool_read_ahead         | 2713        |
| Innodb_buffer_pool_read_ahead_evicted | 0           |
| Innodb_buffer_pool_read_requests      | 33069332    |
| Innodb_buffer_pool_reads              | 558         |
| Innodb_buffer_pool_wait_free          | 0           |
| Innodb_buffer_pool_write_requests     | 11985961    |
+---------------------------------------+-------------+
```

15.14.6 InnoDB INFORMATION_SCHEMA Metrics Table

The `INNODB_METRICS` table stores data for `InnoDB` performance and resource-related counters:

The columns of the `INNODB_METRICS` table are shown in the following example. For a description of each column, see Section 24.31.16, “The INFORMATION_SCHEMA INNODB_METRICS Table”.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts" \G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 46273
      MAX_COUNT: 46273
      MIN_COUNT: NULL
      AVG_COUNT: 492.2659574468085
    COUNT_RESET: 46273
MAX_COUNT_RESET: 46273
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: NULL
   TIME_ENABLED: 2014-11-28 16:07:53
  TIME_DISABLED: NULL
   TIME_ELAPSED: 94
     TIME_RESET: NULL
         STATUS: enabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

Enabling, Disabling, and Resetting Counters

You can enable, disable, and reset counters using the following configuration options:

- `innodb_monitor_enable`: Enables one or more counters.

  ```
  SET GLOBAL innodb_monitor_enable = [counter-name|module_name|pattern|all];
  ```

- `innodb_monitor_disable`: Disables one or more counters.

  ```
  SET GLOBAL innodb_monitor_disable = [counter-name|module_name|pattern|all];
  ```

- `innodb_monitor_reset`: Resets the count value for one or more counters to zero.

  ```
  SET GLOBAL innodb_monitor_reset = [counter-name|module_name|pattern|all];
  ```

- `innodb_monitor_reset_all`: Resets all values for one or more counters. A counter must be disabled before using `innodb_monitor_reset_all`.

  ```
  SET GLOBAL innodb_monitor_reset_all = [counter-name|module_name|pattern|all];
  ```

You can also enable counters and counter modules at startup using the MySQL server configuration file. For example, to enable the `log` module, `metadata_table_handles_opened` and `metadata_table_handles_closed` counters, enter the following line in the `[mysqld]` section of your `my.cnf` configuration file.

```
[mysqld]
innodb_monitor_enable = module_recovery,metadata_table_handles_opened,metadata_table_handles_closed
```

When enabling multiple counters or modules in your configuration file, you must specify the `innodb_monitor_enable` configuration option followed by counter and module names separated by a comma, as shown in the example above. Only the `innodb_monitor_enable` option can be used in your configuration file. The disable and reset configuration options are only supported on the command line.

Note

Because each counter imposes some degree of runtime overhead on the server, typically you enable more counters on test and development servers during experimentation and benchmarking, and only enable counters on production servers to diagnose known issues or monitor aspects that are likely to be bottlenecks for a particular server and workload.

Counters

The counters represented in the `INNODB_METRICS` table are subject to change, so for the most up-to-date list, query a running MySQL server. The list below shows counters that are available as of MySQL 8.0.

Counters that are enabled by default correspond to those used by `SHOW ENGINE INNODB STATUS`. Counters used by `SHOW ENGINE INNODB STATUS` are always “on” at a system level but you can disable these counters for the `INNODB_METRICS` table, as required. Also, counter status is not persistent. Unless specified otherwise, counters revert to their default enabled or disabled status when the server is restarted.

If you run programs that would be affected by additions or changes to the `INNODB_METRICS` table, it is recommended that you review releases notes and query the `INNODB_METRICS` table for the new release prior to upgrading.

```
mysql> SELECT name, subsystem, status FROM INFORMATION_SCHEMA.INNODB_METRICS ORDER BY NAME;
+----------------------------------------+--------------------+----------+
| name                                   | subsystem          | status   |
+----------------------------------------+--------------------+----------+
| adaptive_hash_pages_added              | adaptive_hash_index | disabled |
| adaptive_hash_pages_removed            | adaptive_hash_index | disabled |
| adaptive_hash_rows_added               | adaptive_hash_index | disabled |
| adaptive_hash_rows_deleted_no_hash_entry | adaptive_hash_index | disabled |
| adaptive_hash_rows_removed             | adaptive_hash_index | disabled |
| adaptive_hash_rows_updated             | adaptive_hash_index | disabled |
| adaptive_hash_searches                 | adaptive_hash_index | enabled  |
| adaptive_hash_searches_btree           | adaptive_hash_index | enabled  |
| buffer_data_reads                      | buffer             | enabled  |
| buffer_data_written                    | buffer             | enabled  |
| buffer_flush_adaptive                  | buffer             | disabled |
| buffer_flush_adaptive_avg_pass         | buffer             | disabled |
| buffer_flush_adaptive_avg_time_est     | buffer             | disabled |
| buffer_flush_adaptive_avg_time_slot    | buffer             | disabled |
| buffer_flush_adaptive_avg_time_thread  | buffer             | disabled |
| buffer_flush_adaptive_pages            | buffer             | disabled |
| buffer_flush_adaptive_total_pages      | buffer             | disabled |
| buffer_flush_avg_page_rate             | buffer             | disabled |
| buffer_flush_avg_pass                  | buffer             | disabled |
| buffer_flush_avg_time                  | buffer             | disabled |
| buffer_flush_background                | buffer             | disabled |
| buffer_flush_background_pages          | buffer             | disabled |
| buffer_flush_background_total_pages    | buffer             | disabled |
| buffer_flush_batches                   | buffer             | disabled |
| buffer_flush_batch_num_scan            | buffer             | disabled |
| buffer_flush_batch_pages               | buffer             | disabled |
| buffer_flush_batch_scanned             | buffer             | disabled |
| buffer_flush_batch_scanned_per_call    | buffer             | disabled |
| buffer_flush_batch_total_pages         | buffer             | disabled |
| buffer_flush_lsn_avg_rate              | buffer             | disabled |
```

```
| buffer_flush_neighbor                    | buffer             | disabled |
| buffer_flush_neighbor_pages              | buffer             | disabled |
| buffer_flush_neighbor_total_pages        | buffer             | disabled |
| buffer_flush_n_to_flush_by_age           | buffer             | disabled |
| buffer_flush_n_to_flush_requested        | buffer             | disabled |
| buffer_flush_pct_for_dirty               | buffer             | disabled |
| buffer_flush_pct_for_lsn                 | buffer             | disabled |
| buffer_flush_sync                        | buffer             | disabled |
| buffer_flush_sync_pages                  | buffer             | disabled |
| buffer_flush_sync_total_pages            | buffer             | disabled |
| buffer_flush_sync_waits                  | buffer             | disabled |
| buffer_LRU_batches_evict                 | buffer             | disabled |
| buffer_LRU_batches_flush                 | buffer             | disabled |
| buffer_LRU_batch_evict_pages             | buffer             | disabled |
| buffer_LRU_batch_evict_total_pages       | buffer             | disabled |
| buffer_LRU_batch_flush_avg_pass          | buffer             | disabled |
| buffer_LRU_batch_flush_avg_time_est      | buffer             | disabled |
| buffer_LRU_batch_flush_avg_time_slot     | buffer             | disabled |
| buffer_LRU_batch_flush_avg_time_thread   | buffer             | disabled |
| buffer_LRU_batch_flush_pages             | buffer             | disabled |
| buffer_LRU_batch_flush_total_pages       | buffer             | disabled |
| buffer_LRU_batch_num_scan                | buffer             | disabled |
| buffer_LRU_batch_scanned                 | buffer             | disabled |
| buffer_LRU_batch_scanned_per_call        | buffer             | disabled |
| buffer_LRU_get_free_loops                | buffer             | disabled |
| buffer_LRU_get_free_search               | Buffer             | disabled |
| buffer_LRU_get_free_waits                | buffer             | disabled |
| buffer_LRU_search_num_scan               | buffer             | disabled |
| buffer_LRU_search_scanned                | buffer             | disabled |
| buffer_LRU_search_scanned_per_call       | buffer             | disabled |
| buffer_LRU_single_flush_failure_count    | Buffer             | disabled |
| buffer_LRU_single_flush_num_scan         | buffer             | disabled |
| buffer_LRU_single_flush_scanned          | buffer             | disabled |
| buffer_LRU_single_flush_scanned_per_call | buffer             | disabled |
| buffer_LRU_unzip_search_num_scan         | buffer             | disabled |
| buffer_LRU_unzip_search_scanned          | buffer             | disabled |
| buffer_LRU_unzip_search_scanned_per_call | buffer             | disabled |
| buffer_pages_created                     | buffer             | enabled  |
| buffer_pages_read                        | buffer             | enabled  |
| buffer_pages_written                     | buffer             | enabled  |
| buffer_page_read_blob                    | buffer_page_io     | disabled |
| buffer_page_read_fsp_hdr                 | buffer_page_io     | disabled |
| buffer_page_read_ibuf_bitmap             | buffer_page_io     | disabled |
| buffer_page_read_ibuf_free_list          | buffer_page_io     | disabled |
| buffer_page_read_index_ibuf_leaf         | buffer_page_io     | disabled |
| buffer_page_read_index_ibuf_non_leaf     | buffer_page_io     | disabled |
| buffer_page_read_index_inode             | buffer_page_io     | disabled |
| buffer_page_read_index_leaf              | buffer_page_io     | disabled |
| buffer_page_read_index_non_leaf          | buffer_page_io     | disabled |
| buffer_page_read_other                   | buffer_page_io     | disabled |
| buffer_page_read_system_page             | buffer_page_io     | disabled |
| buffer_page_read_trx_system              | buffer_page_io     | disabled |
| buffer_page_read_undo_log                | buffer_page_io     | disabled |
| buffer_page_read_xdes                    | buffer_page_io     | disabled |
| buffer_page_read_zblob                   | buffer_page_io     | disabled |
| buffer_page_read_zblob2                  | buffer_page_io     | disabled |
| buffer_page_written_blob                 | buffer_page_io     | disabled |
| buffer_page_written_fsp_hdr              | buffer_page_io     | disabled |
| buffer_page_written_ibuf_bitmap          | buffer_page_io     | disabled |
| buffer_page_written_ibuf_free_list       | buffer_page_io     | disabled |
| buffer_page_written_index_ibuf_leaf      | buffer_page_io     | disabled |
| buffer_page_written_index_ibuf_non_leaf  | buffer_page_io     | disabled |
| buffer_page_written_index_inode          | buffer_page_io     | disabled |
| buffer_page_written_index_leaf           | buffer_page_io     | disabled |
| buffer_page_written_index_non_leaf       | buffer_page_io     | disabled |
| buffer_page_written_other                | buffer_page_io     | disabled |
| buffer_page_written_system_page          | buffer_page_io     | disabled |
| buffer_page_written_trx_system           | buffer_page_io     | disabled |
| buffer_page_written_undo_log             | buffer_page_io     | disabled |
| buffer_page_written_xdes                 | buffer_page_io     | disabled |
| buffer_page_written_zblob                | buffer_page_io     | disabled |
| buffer_page_written_zblob2               | buffer_page_io     | disabled |
```

```
| buffer_pool_bytes_data                | buffer             | enabled  |
| buffer_pool_bytes_dirty               | buffer             | enabled  |
| buffer_pool_pages_data                | buffer             | enabled  |
| buffer_pool_pages_dirty               | buffer             | enabled  |
| buffer_pool_pages_free                | buffer             | enabled  |
| buffer_pool_pages_misc                | buffer             | enabled  |
| buffer_pool_pages_total               | buffer             | enabled  |
| buffer_pool_reads                     | buffer             | enabled  |
| buffer_pool_read_ahead                | buffer             | enabled  |
| buffer_pool_read_ahead_evicted        | buffer             | enabled  |
| buffer_pool_read_requests             | buffer             | enabled  |
| buffer_pool_size                      | server             | enabled  |
| buffer_pool_wait_free                 | buffer             | enabled  |
| buffer_pool_write_requests            | buffer             | enabled  |
| compression_pad_decrements            | compression        | disabled |
| compression_pad_increments            | compression        | disabled |
| compress_pages_compressed             | compression        | disabled |
| compress_pages_decompressed           | compression        | disabled |
| ddl_background_drop_indexes           | ddl                | disabled |
| ddl_background_drop_tables            | ddl                | disabled |
| ddl_log_file_alter_table              | ddl                | disabled |
| ddl_online_create_index               | ddl                | disabled |
| ddl_pending_alter_table               | ddl                | disabled |
| ddl_sort_file_alter_table             | ddl                | disabled |
| dml_deletes                           | dml                | enabled  |
| dml_inserts                           | dml                | enabled  |
| dml_reads                             | dml                | disabled |
| dml_updates                           | dml                | enabled  |
| file_num_open_files                   | file_system        | enabled  |
| ibuf_merges                           | change_buffer      | enabled  |
| ibuf_merges_delete                    | change_buffer      | enabled  |
| ibuf_merges_delete_mark               | change_buffer      | enabled  |
| ibuf_merges_discard_delete            | change_buffer      | enabled  |
| ibuf_merges_discard_delete_mark       | change_buffer      | enabled  |
| ibuf_merges_discard_insert            | change_buffer      | enabled  |
| ibuf_merges_insert                    | change_buffer      | enabled  |
| ibuf_size                             | change_buffer      | enabled  |
| icp_attempts                          | icp                | disabled |
| icp_match                             | icp                | disabled |
| icp_no_match                          | icp                | disabled |
| icp_out_of_range                      | icp                | disabled |
| index_page_discards                   | index              | disabled |
| index_page_merge_attempts             | index              | disabled |
| index_page_merge_successful           | index              | disabled |
| index_page_reorg_attempts             | index              | disabled |
| index_page_reorg_successful           | index              | disabled |
| index_page_splits                     | index              | disabled |
| innodb_activity_count                 | server             | enabled  |
| innodb_background_drop_table_usec     | server             | disabled |
| innodb_checkpoint_usec                | server             | disabled |
| innodb_dblwr_pages_written            | server             | enabled  |
| innodb_dblwr_writes                   | server             | enabled  |
| innodb_dict_lru_count                 | server             | disabled |
| innodb_dict_lru_usec                  | server             | disabled |
| innodb_ibuf_merge_usec                | server             | disabled |
| innodb_log_flush_usec                 | server             | disabled |
| innodb_master_active_loops            | server             | disabled |
| innodb_master_idle_loops              | server             | disabled |
| innodb_master_purge_usec              | server             | disabled |
| innodb_master_thread_sleeps           | server             | disabled |
| innodb_mem_validate_usec              | server             | disabled |
| innodb_page_size                      | server             | enabled  |
| innodb_rwlock_sx_os_waits             | server             | enabled  |
| innodb_rwlock_sx_spin_rounds          | server             | enabled  |
| innodb_rwlock_sx_spin_waits           | server             | enabled  |
| innodb_rwlock_s_os_waits              | server             | enabled  |
| innodb_rwlock_s_spin_rounds           | server             | enabled  |
| innodb_rwlock_s_spin_waits            | server             | enabled  |
| innodb_rwlock_x_os_waits              | server             | enabled  |
| innodb_rwlock_x_spin_rounds           | server             | enabled  |
| innodb_rwlock_x_spin_waits            | server             | enabled  |
| lock_deadlocks                        | lock               | enabled  |
```

```
| lock_rec_locks                          | lock               | disabled |
| lock_rec_lock_created                   | lock               | disabled |
| lock_rec_lock_removed                   | lock               | disabled |
| lock_rec_lock_requests                  | lock               | disabled |
| lock_rec_lock_waits                     | lock               | disabled |
| lock_row_lock_current_waits             | lock               | enabled  |
| lock_row_lock_time                      | lock               | enabled  |
| lock_row_lock_time_avg                  | lock               | enabled  |
| lock_row_lock_time_max                  | lock               | enabled  |
| lock_row_lock_waits                     | lock               | enabled  |
| lock_table_locks                        | lock               | disabled |
| lock_table_lock_created                 | lock               | disabled |
| lock_table_lock_removed                 | lock               | disabled |
| lock_table_lock_waits                   | lock               | disabled |
| lock_timeouts                           | lock               | enabled  |
| log_checkpoints                         | recovery           | disabled |
| log_lsn_buf_pool_oldest                 | recovery           | disabled |
| log_lsn_checkpoint_age                  | recovery           | disabled |
| log_lsn_current                         | recovery           | disabled |
| log_lsn_last_checkpoint                 | recovery           | disabled |
| log_lsn_last_flush                      | recovery           | disabled |
| log_max_modified_age_async              | recovery           | disabled |
| log_max_modified_age_sync               | recovery           | disabled |
| log_num_log_io                          | recovery           | disabled |
| log_padded                              | recovery           | enabled  |
| log_pending_checkpoint_writes           | recovery           | disabled |
| log_pending_log_flushes                 | recovery           | disabled |
| log_waits                               | recovery           | enabled  |
| log_writes                              | recovery           | enabled  |
| log_write_requests                      | recovery           | enabled  |
| metadata_table_handles_closed           | metadata           | disabled |
| metadata_table_handles_opened           | metadata           | disabled |
| metadata_table_reference_count          | metadata           | disabled |
| os_data_fsyncs                          | os                 | enabled  |
| os_data_reads                           | os                 | enabled  |
| os_data_writes                          | os                 | enabled  |
| os_log_bytes_written                    | os                 | enabled  |
| os_log_fsyncs                           | os                 | enabled  |
| os_log_pending_fsyncs                   | os                 | enabled  |
| os_log_pending_writes                   | os                 | enabled  |
| os_pending_reads                        | os                 | disabled |
| os_pending_writes                       | os                 | disabled |
| purge_del_mark_records                  | purge              | disabled |
| purge_dml_delay_usec                    | purge              | disabled |
| purge_invoked                           | purge              | disabled |
| purge_resume_count                      | purge              | disabled |
| purge_stop_count                        | purge              | disabled |
| purge_undo_log_pages                    | purge              | disabled |
| purge_upd_exist_or_extern_records       | purge              | disabled |
| trx_active_transactions                 | transaction        | disabled |
| trx_commits_insert_update               | transaction        | disabled |
| trx_nl_ro_commits                       | transaction        | disabled |
| trx_rollbacks                           | transaction        | disabled |
| trx_rollbacks_savepoint                 | transaction        | disabled |
| trx_rollback_active                     | transaction        | disabled |
| trx_ro_commits                          | transaction        | disabled |
| trx_rseg_current_size                   | transaction        | disabled |
| trx_rseg_history_len                    | transaction        | enabled  |
| trx_rw_commits                          | transaction        | disabled |
| trx_undo_slots_cached                   | transaction        | disabled |
| trx_undo_slots_used                     | transaction        | disabled |
+-----------------------------------------+--------------------+----------+
235 rows in set (0.01 sec)
```

Counter Modules

The module names correspond to, but are not identical to, the values from the `SUBSYSTEM` column of the `INNODB_METRICS` table. Rather enabling, disabling, or resetting counters individually, you can use module names to quickly enable, disable, or reset all counters for a particular subsystem. For example, use `module_dml` to enable all counters associated with the `dml` subsystem.

```
mysql> SET GLOBAL innodb_monitor_enable = module_dml;

mysql> SELECT name, subsystem, status FROM INFORMATION_SCHEMA.INNODB_METRICS
WHERE subsystem ='dml';
+-------------+-----------+---------+
| name        | subsystem | status  |
+-------------+-----------+---------+
| dml_reads   | dml       | enabled |
| dml_inserts | dml       | enabled |
| dml_deletes | dml       | enabled |
| dml_updates | dml       | enabled |
+-------------+-----------+---------+
```

Here are the values you can use for *`module_name`* with the `innodb_monitor_enable` and related configuration options, along with the corresponding `SUBSYSTEM` names:

- `module_adaptive_hash` (subsystem = `adaptive_hash_index`)
- `module_buffer` (subsystem = `buffer`)
- `module_buffer_page` (subsystem = `buffer_page_io`)
- `module_compress` (subsystem = `compression`)
- `module_ddl` (subsystem = `ddl`)
- `module_dml` (subsystem = `dml`)
- `module_file` (subsystem = `file_system`)
- `module_ibuf_system` (subsystem = `change_buffer`)
- `module_icp` (subsystem = `icp`)
- `module_index` (subsystem = `index`)
- `module_innodb` (subsystem = `innodb`)
- `module_lock` (subsystem = `lock`)
- `module_log` (subsystem = `recovery`)
- `module_metadata` (subsystem = `metadata`)
- `module_os` (subsystem = `os`)
- `module_purge` (subsystem = `purge`)
- `module_trx` (subsystem = `transaction`)

Example 15.11 Working with INNODB_METRICS Table Counters

This example demonstrates enabling, disabling, and resetting a counter, and querying counter data in the `INNODB_METRICS` table.

1. Create a simple `InnoDB` table:

   ```
   mysql> USE test;
   Database changed

   mysql> CREATE TABLE t1 (c1 INT) ENGINE=INNODB;
   Query OK, 0 rows affected (0.02 sec)
   ```

2. Enable the `dml_inserts` counter.

   ```
   mysql> SET GLOBAL innodb_monitor_enable = dml_inserts;
   Query OK, 0 rows affected (0.01 sec)
   ```

A description of the `dml_inserts` counter can be found in the `COMMENT` column of the `INNODB_METRICS` table:

```
mysql> SELECT NAME, COMMENT FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts";
+-------------+------------------------+
| NAME        | COMMENT                |
+-------------+------------------------+
| dml_inserts | Number of rows inserted |
+-------------+------------------------+
```

3. Query the `INNODB_METRICS` table for the `dml_inserts` counter data. Because no DML operations have been performed, the counter values are zero or NULL. The `TIME_ENABLED` and `TIME_ELAPSED` values indicate when the counter was last enabled and how many seconds have elapsed since this time.

```
mysql>  SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts" \G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 0
      MAX_COUNT: 0
      MIN_COUNT: NULL
      AVG_COUNT: 0
    COUNT_RESET: 0
MAX_COUNT_RESET: 0
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: NULL
   TIME_ENABLED: 2014-12-04 14:18:28
  TIME_DISABLED: NULL
   TIME_ELAPSED: 28
     TIME_RESET: NULL
         STATUS: enabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

4. Insert three rows of data into the table.

```
mysql> INSERT INTO t1 values(1);
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO t1 values(2);
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO t1 values(3);
Query OK, 1 row affected (0.00 sec)
```

5. Query the `INNODB_METRICS` table again for the `dml_inserts` counter data. A number of counter values have now incremented including `COUNT`, `MAX_COUNT`, `AVG_COUNT`, and `COUNT_RESET`. Refer to the `INNODB_METRICS` table definition for descriptions of these values.

```
mysql>  SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts"\G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 3
      MAX_COUNT: 3
      MIN_COUNT: NULL
      AVG_COUNT: 0.046153846153846156
    COUNT_RESET: 3
MAX_COUNT_RESET: 3
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: NULL
   TIME_ENABLED: 2014-12-04 14:18:28
  TIME_DISABLED: NULL
   TIME_ELAPSED: 65
     TIME_RESET: NULL
         STATUS: enabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

6. Reset the `dml_inserts` counter, and query the `INNODB_METRICS` table again for the `dml_inserts` counter data. The `%_RESET` values that were reported previously, such as `COUNT_RESET` and `MAX_RESET`, are set back to zero. Values such as `COUNT`, `MAX_COUNT`, and `AVG_COUNT`, which cumulatively collect data from the time the counter is enabled, are unaffected by the reset.

```
mysql> SET GLOBAL innodb_monitor_reset = dml_inserts;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts"\G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 3
      MAX_COUNT: 3
      MIN_COUNT: NULL
      AVG_COUNT: 0.03529411764705882
    COUNT_RESET: 0
MAX_COUNT_RESET: 0
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: 0
   TIME_ENABLED: 2014-12-04 14:18:28
  TIME_DISABLED: NULL
   TIME_ELAPSED: 85
     TIME_RESET: 2014-12-04 14:19:44
         STATUS: enabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

7. To reset all counter values, you must first disable the counter. Disabling the counter sets the `STATUS` value to `disbaled`.

```
mysql> SET GLOBAL innodb_monitor_disable = dml_inserts;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts"\G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 3
      MAX_COUNT: 3
      MIN_COUNT: NULL
      AVG_COUNT: 0.030612244897959183
    COUNT_RESET: 0
MAX_COUNT_RESET: 0
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: 0
   TIME_ENABLED: 2014-12-04 14:18:28
  TIME_DISABLED: 2014-12-04 14:20:06
   TIME_ELAPSED: 98
     TIME_RESET: NULL
         STATUS: disabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

Note

Wildcard match is supported for counter and module names. For example, instead of specifying the full `dml_inserts` counter name, you can specify `dml_i%`. You can also enable, disable, or reset multiple counters or modules at once using a wildcard match. For example, specify `dml_%` to enable, disable, or reset all counters that begin with `dml_%`.

8. After the counter is disabled, you can reset all counter values using the `innodb_monitor_reset_all` option. All values are set to zero or NULL.

```
mysql> SET GLOBAL innodb_monitor_reset_all = dml_inserts;
Query OK, 0 rows affected (0.00 sec)
```

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts"\G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 0
      MAX_COUNT: NULL
      MIN_COUNT: NULL
      AVG_COUNT: NULL
    COUNT_RESET: 0
MAX_COUNT_RESET: NULL
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: NULL
   TIME_ENABLED: NULL
  TIME_DISABLED: NULL
   TIME_ELAPSED: NULL
     TIME_RESET: NULL
         STATUS: disabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

15.14.7 InnoDB INFORMATION_SCHEMA Temporary Table Information Table

The `INNODB_TEMP_TABLE_INFO` table provides users with a snapshot of active `InnoDB` temporary tables. The table contains metadata about all user and system-created temporary tables that are active within a given `InnoDB` instance with the exception of optimized temporary tables that are used internally by `InnoDB`.

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA LIKE 'INNODB_TEMP%';
+--------------------------------------------+
| Tables_in_INFORMATION_SCHEMA (INNODB_TEMP%) |
+--------------------------------------------+
| INNODB_TEMP_TABLE_INFO                     |
+--------------------------------------------+
```

For the table definition, see Section 24.31.27, “The INFORMATION_SCHEMA INNODB_TEMP_TABLE_INFO Table”.

Example 15.12 INNODB_TEMP_TABLE_INFO

This example demonstrates characteristics of the `INNODB_TEMP_TABLE_INFO` table.

1. Create a simple `InnoDB` temporary table with a single column:

   ```
   mysql> CREATE TEMPORARY TABLE t1 (c1 INT PRIMARY KEY) ENGINE=INNODB;
   Query OK, 0 rows affected (0.00 sec)
   ```

2. Query the `INNODB_TEMP_TABLE_INFO` table to view the temporary table's metadata.

   ```
   mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_TEMP_TABLE_INFO\G
   *************************** 1. row ***************************
              TABLE_ID: 194
                  NAME: #sql7a79_1_0
               N_COLS: 4
                SPACE: 182
   ```

   The `TABLE_ID` is a unique identifier for the temporary table. The `NAME` column displays the system-generated name for the temporary table, which is prefixed with “#sql”. The number of columns (`N_COLS`) is 4 rather than 1 because `InnoDB` always creates three hidden table columns (`DB_ROW_ID`, `DB_TRX_ID`, and `DB_ROLL_PTR`).

3. Restart MySQL and query the `INNODB_TEMP_TABLE_INFO` table.

   ```
   mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_TEMP_TABLE_INFO\G
   Empty set (0.00 sec)
   ```

   An empty set is returned because the `INNODB_TEMP_TABLE_INFO` table and the data within it are not persisted to disk on server shutdown.

4. Create a new temporary table.

```
mysql> CREATE TEMPORARY TABLE t1 (c1 INT PRIMARY KEY) ENGINE=INNODB;
Query OK, 0 rows affected (0.00 sec)
```

5. Query the INNODB_TEMP_TABLE_INFO table to view the temporary table's metadata.

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_TEMP_TABLE_INFO\G
*************************** 1. row ***************************
            TABLE_ID: 196
                NAME: #sql7b0e_1_0
              N_COLS: 4
               SPACE: 184
```

The SPACE ID is new because it is dynamically generated on server restart.

15.14.8 Retrieving InnoDB Tablespace Metadata from INFORMATION_SCHEMA.FILES

The INFORMATION_SCHEMA.FILES table provides metadata about all InnoDB tablespace types including file-per-table tablespaces, general tablespaces, the system tablespace, temporary table tablespaces, and undo tablespaces (if present).

This section provides InnoDB-specific usage examples. For more information about data provided by the INFORMATION_SCHEMA.FILES table, see Section 24.9, “The INFORMATION_SCHEMA FILES Table”.

Note

The INNODB_SYS_TABLESPACES and INNODB_SYS_DATAFILES tables also provide metadata about InnoDB tablespaces, but data is limited to file-per-table and general tablespaces.

This query retrieves metadata about the InnoDB system tablespace from fields of the INFORMATION_SCHEMA.FILES table that are pertinent to InnoDB tablespaces. INFORMATION_SCHEMA.FILES fields that are not relevant to InnoDB always return NULL, and are excluded from the query.

```
mysql> SELECT FILE_ID, FILE_NAME, FILE_TYPE, TABLESPACE_NAME, FREE_EXTENTS,
TOTAL_EXTENTS,  EXTENT_SIZE, INITIAL_SIZE, MAXIMUM_SIZE, AUTOEXTEND_SIZE, DATA_FREE, STATUS ENGINE
FROM INFORMATION_SCHEMA.FILES WHERE TABLESPACE_NAME LIKE 'innodb_system' \G
*************************** 1. row ***************************
        FILE_ID: 0
      FILE_NAME: ./ibdata1
      FILE_TYPE: TABLESPACE
TABLESPACE_NAME: innodb_system
   FREE_EXTENTS: 0
  TOTAL_EXTENTS: 12
    EXTENT_SIZE: 1048576
   INITIAL_SIZE: 12582912
   MAXIMUM_SIZE: NULL
AUTOEXTEND_SIZE: 67108864
      DATA_FREE: 4194304
         ENGINE: NORMAL
```

This query retrieves the FILE_ID (equivalent to the space ID) and the FILE_NAME (which includes path information) for InnoDB file-per-table and general tablespaces. File-per-table and general tablespaces have a .ibd file extension.

```
mysql> SELECT FILE_ID, FILE_NAME FROM INFORMATION_SCHEMA.FILES
WHERE FILE_NAME LIKE '%.ibd%' ORDER BY FILE_ID;
    +---------+--------------------------------------+
    | FILE_ID | FILE_NAME                            |
    +---------+--------------------------------------+
    |       2 | ./mysql/plugin.ibd                   |
    |       3 | ./mysql/servers.ibd                  |
    |       4 | ./mysql/help_topic.ibd               |
```

```
|       5 | ./mysql/help_category.ibd             |
|       6 | ./mysql/help_relation.ibd             |
|       7 | ./mysql/help_keyword.ibd              |
|       8 | ./mysql/time_zone_name.ibd            |
|       9 | ./mysql/time_zone.ibd                 |
|      10 | ./mysql/time_zone_transition.ibd      |
|      11 | ./mysql/time_zone_transition_type.ibd |
|      12 | ./mysql/time_zone_leap_second.ibd     |
|      13 | ./mysql/innodb_table_stats.ibd        |
|      14 | ./mysql/innodb_index_stats.ibd        |
|      15 | ./mysql/slave_relay_log_info.ibd      |
|      16 | ./mysql/slave_master_info.ibd         |
|      17 | ./mysql/slave_worker_info.ibd         |
|      18 | ./mysql/gtid_executed.ibd             |
|      19 | ./mysql/server_cost.ibd               |
|      20 | ./mysql/engine_cost.ibd               |
|      21 | ./sys/sys_config.ibd                  |
|      23 | ./test/t1.ibd                         |
|      26 | /home/user/test/test/t2.ibd           |
+---------+---------------------------------------+
```

This query retrieves the `FILE_ID` and `FILE_NAME` for `InnoDB` temporary tablespaces. Temporary tablespace file names are prefixed by `ibtmp`.

```
mysql> SELECT FILE_ID, FILE_NAME FROM INFORMATION_SCHEMA.FILES
WHERE FILE_NAME LIKE '%ibtmp%';
+---------+-----------+
| FILE_ID | FILE_NAME |
+---------+-----------+
|      22 | ./ibtmp1  |
+---------+-----------+
```

Similarly, `InnoDB` undo tablespace file names are prefixed by `undo`. The following query returns the `FILE_ID` and `FILE_NAME` for `InnoDB` undo tablespaces, if separate undo tablespaces are configured.

```
mysql> SELECT FILE_ID, FILE_NAME FROM INFORMATION_SCHEMA.FILES
WHERE FILE_NAME LIKE '%undo%';
```

15.15 InnoDB Integration with MySQL Performance Schema

This section provides a brief introduction to `InnoDB` integration with Performance Schema. For comprehensive Performance Schema documentation, see Chapter 25, *MySQL Performance Schema*.

You can profile certain internal `InnoDB` operations using the MySQL Performance Schema feature. This type of tuning is primarily for expert users who evaluate optimization strategies to overcome performance bottlenecks. DBAs can also use this feature for capacity planning, to see whether their typical workload encounters any performance bottlenecks with a particular combination of CPU, RAM, and disk storage; and if so, to judge whether performance can be improved by increasing the capacity of some part of the system.

To use this feature to examine `InnoDB` performance:

- You must be generally familiar with how to use the Performance Schema feature. For example, you should know how enable instruments and consumers, and how to query `performance_schema` tables to retrieve data. For an introductory overview, see Section 25.1, “Performance Schema Quick Start”.
- You should be familiar with Performance Schema instruments that are available for `InnoDB`. To view `InnoDB`-related instruments, you can query the `setup_instruments` table for instrument names that contain '`innodb`'.

```
mysql> SELECT * FROM setup_instruments WHERE NAME LIKE '%innodb%';
+---------------------------------------------------------+---------+-------+
| NAME                                                    | ENABLED | TIMED |
+---------------------------------------------------------+---------+-------+
| wait/synch/mutex/innodb/commit_cond_mutex               | NO      | NO    |
| wait/synch/mutex/innodb/innobase_share_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/autoinc_mutex                   | NO      | NO    |
```

```
| wait/synch/mutex/innodb/buf_pool_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_zip_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/cache_last_read_mutex         | NO      | NO    |
| wait/synch/mutex/innodb/dict_foreign_err_mutex        | NO      | NO    |
| wait/synch/mutex/innodb/dict_sys_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/recalc_pool_mutex             | NO      | NO    |
...
| wait/io/file/innodb/innodb_data_file                  | YES     | YES   |
| wait/io/file/innodb/innodb_log_file                   | YES     | YES   |
| wait/io/file/innodb/innodb_temp_file                  | YES     | YES   |
| stage/innodb/alter table (end)                        | YES     | YES   |
| stage/innodb/alter table (flush)                      | YES     | YES   |
| stage/innodb/alter table (insert)                     | YES     | YES   |
| stage/innodb/alter table (log apply index)            | YES     | YES   |
| stage/innodb/alter table (log apply table)            | YES     | YES   |
| stage/innodb/alter table (merge sort)                 | YES     | YES   |
| stage/innodb/alter table (read PK and internal sort)  | YES     | YES   |
| stage/innodb/buffer pool load                         | YES     | YES   |
| memory/innodb/buf_buf_pool                            | NO      | NO    |
| memory/innodb/dict_stats_bg_recalc_pool_t             | NO      | NO    |
| memory/innodb/dict_stats_index_map_t                  | NO      | NO    |
| memory/innodb/dict_stats_n_diff_on_level              | NO      | NO    |
| memory/innodb/other                                   | NO      | NO    |
| memory/innodb/row_log_buf                             | NO      | NO    |
| memory/innodb/row_merge_sort                          | NO      | NO    |
| memory/innodb/std                                     | NO      | NO    |
| memory/innodb/sync_debug_latches                      | NO      | NO    |
| memory/innodb/trx_sys_t::rw_trx_ids                   | NO      | NO    |
...
+-------------------------------------------------------+---------+-------+
155 rows in set (0.00 sec)
```

For additional information about the instrumented `InnoDB` objects, you can query Performance Schema instances tables, which provide additional information about instrumented objects. Instance tables relevant to `InnoDB` include:

- The `mutex_instances` table
- The `rwlock_instances` table
- The `cond_instances` table
- The `file_instances` table

Note

Mutexes and RW-locks related to the `InnoDB` buffer pool are not included in this coverage; the same applies to the output of the `SHOW ENGINE INNODB MUTEX` command.

For example, to view information about instrumented `InnoDB` file objects seen by the Performance Schema when executing file I/O instrumentation, you might issue the following query:

```
mysql> SELECT * FROM file_instances WHERE EVENT_NAME LIKE '%innodb%'\G
*************************** 1. row ***************************
 FILE_NAME: /path/to/mysql-8.0/data/ibdata1
EVENT_NAME: wait/io/file/innodb/innodb_data_file
OPEN_COUNT: 3
*************************** 2. row ***************************
 FILE_NAME: /path/to/mysql-8.0/data/ib_logfile0
EVENT_NAME: wait/io/file/innodb/innodb_log_file
OPEN_COUNT: 2
*************************** 3. row ***************************
 FILE_NAME: /path/to/mysql-8.0/data/ib_logfile1
EVENT_NAME: wait/io/file/innodb/innodb_log_file
OPEN_COUNT: 2
*************************** 4. row ***************************
 FILE_NAME: /path/to/mysql-8.0/data/mysql/engine_cost.ibd
EVENT_NAME: wait/io/file/innodb/innodb_data_file
```

```
OPEN_COUNT: 3
...
```

- You should be familiar with `performance_schema` tables that store `InnoDB` event data. Tables relevant to `InnoDB`-related events include:

  - The Wait Event tables, which store wait events.

  - The Summary tables, which provide aggregated information for terminated events over time. Summary tables include file I/O summary tables, which aggregate information about I/O operations.

  - Stage Event tables, which store event data for `InnoDB ALTER TABLE` and buffer pool load operations. For more information, see Section 15.15.1, “Monitoring ALTER TABLE Progress for InnoDB Tables Using Performance Schema”, and Monitoring Buffer Pool Load Progress Using Performance Schema.

  If you are only interested in `InnoDB`-related objects, use the clause `WHERE EVENT_NAME LIKE '%innodb%'` or `WHERE NAME LIKE '%innodb%'` (as required) when querying these tables.

15.15.1 Monitoring ALTER TABLE Progress for InnoDB Tables Using Performance Schema

You can monitor `ALTER TABLE` progress for `InnoDB` tables using Performance Schema.

There are seven stage events that represent different phases of `ALTER TABLE`. Each stage event reports a running total of `WORK_COMPLETED` and `WORK_ESTIMATED` for the overall `ALTER TABLE` operation as it progresses through its different phases. `WORK_ESTIMATED` is calculated using a formula that takes into account all of the work that `ALTER TABLE` performs, and may be revised during `ALTER TABLE` processing. `WORK_COMPLETED` and `WORK_ESTIMATED` values are an abstract representation of all of the work performed by `ALTER TABLE`.

In order of occurrence, `ALTER TABLE` stage events include:

- `stage/innodb/alter table (read PK and internal sort)`: This stage is active when `ALTER TABLE` is in the reading-primary-key phase. It starts with `WORK_COMPLETED=0` and `WORK_ESTIMATED` set to the estimated number of pages in the primary key. When the stage is completed, `WORK_ESTIMATED` is updated to the actual number of pages in the primary key.

- `stage/innodb/alter table (merge sort)`: This stage is repeated for each index added by the `ALTER TABLE` operation.

- `stage/innodb/alter table (insert)`: This stage is repeated for each index added by the `ALTER TABLE` operation.

- `stage/innodb/alter table (log apply index)`: This stage includes the application of DML log generated while `ALTER TABLE` was running.

- `stage/innodb/alter table (flush)`: Before this stage begins, `WORK_ESTIMATED` is updated with a more accurate estimate, based on the length of the flush list.

- `stage/innodb/alter table (log apply table)`: This stage includes the application of concurrent DML log generated while `ALTER TABLE` was running. The duration of this phase depends on the extent of table changes. This phase is instant if no concurrent DML was run on the table.

- `stage/innodb/alter table (end)`: Includes any remaining work that appeared after the flush phase, such as reapplying DML that was executed on the table while `ALTER TABLE` was running.

Note

`InnoDB ALTER TABLE` stage events do not currently account for the addition of spatial indexes.

ALTER TABLE Monitoring Example Using Performance Schema

The following example demonstrates how to enable the `stage/innodb/alter table%` stage event instruments and related consumer tables to monitor `ALTER TABLE` progress. For information about Performance Schema stage event instruments and related consumers, see Section 25.11.5, "Performance Schema Stage Event Tables".

1. Enable the `stage/innodb/alter%` instruments:

```
mysql> UPDATE setup_instruments SET ENABLED = 'YES' WHERE NAME LIKE 'stage/innodb/alter%';
Query OK, 7 rows affected (0.00 sec)
Rows matched: 7  Changed: 7  Warnings: 0
```

2. Enable the stage event consumer tables, which include `events_stages_current`, `events_stages_history`, and `events_stages_history_long`.

```
mysql> UPDATE setup_consumers SET ENABLED = 'YES' WHERE NAME LIKE '%stages%';
Query OK, 3 rows affected (0.00 sec)
Rows matched: 3  Changed: 3  Warnings: 0
```

3. Run an `ALTER TABLE` operation. In this example, a `middle_name` column is added to the employees table of the employees sample database.

```
mysql> ALTER TABLE employees.employees ADD COLUMN middle_name varchar(14) AFTER first_name;
Query OK, 0 rows affected (9.27 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

4. Check the progress of the `ALTER TABLE` operation by querying the Performance Schema `events_stages_current` table. The stage event shown differs depending on which `ALTER TABLE` phase is currently in progress. The `WORK_COMPLETED` column shows the work completed. The `WORK_ESTIMATED` column provides an estimate of the remaining work.

```
mysql> SELECT EVENT_NAME, WORK_COMPLETED, WORK_ESTIMATED FROM events_stages_current;
+----------------------------------------------------+----------------+----------------+
| EVENT_NAME                                         | WORK_COMPLETED | WORK_ESTIMATED |
+----------------------------------------------------+----------------+----------------+
| stage/innodb/alter table (read PK and internal sort) |            280 |           1245 |
+----------------------------------------------------+----------------+----------------+
1 row in set (0.01 sec)
```

   The `events_stages_current` table returns an empty set if the `ALTER TABLE` operation has completed. In this case, you can check the `events_stages_history` table to view event data for the completed operation. For example:

```
mysql> SELECT EVENT_NAME, WORK_COMPLETED, WORK_ESTIMATED FROM events_stages_history;
+----------------------------------------------------+----------------+----------------+
| EVENT_NAME                                         | WORK_COMPLETED | WORK_ESTIMATED |
+----------------------------------------------------+----------------+----------------+
| stage/innodb/alter table (read PK and internal sort) |            886 |           1213 |
| stage/innodb/alter table (flush)                   |           1213 |           1213 |
| stage/innodb/alter table (log apply table)         |           1597 |           1597 |
| stage/innodb/alter table (end)                     |           1597 |           1597 |
| stage/innodb/alter table (log apply table)         |           1981 |           1981 |
+----------------------------------------------------+----------------+----------------+
5 rows in set (0.00 sec)
```

   As shown above, the `WORK_ESTIMATED` value was revised during `ALTER TABLE` processing. The estimated work after completion of the initial stage is 1213. When `ALTER TABLE` processing completed, `WORK_ESTIMATED` was set to the actual value, which is 1981.

15.15.2 Monitoring InnoDB Mutex Waits Using Performance Schema

A mutex is a synchronization mechanism used in the code to enforce that only one thread at a given time can have access to a common resource. When two or more threads executing in the server need to access the same resource, the threads compete against each other. The first thread to obtain a lock on the mutex causes the other threads to wait until the lock is released.

For `InnoDB` mutexes that are instrumented, mutex waits can be monitored using Performance Schema. Wait event data collected in Performance Schema tables can help identify mutexes with the most waits or the greatest total wait time, for example.

The following example demonstrates how to enable `InnoDB` mutex wait instruments, how to enable associated consumers, and how to query wait event data.

1. To view available `InnoDB` mutex wait instruments, query the Performance Schema `setup_instruments` table. All `InnoDB` mutex wait instruments are disabled by default.

```
mysql> SELECT * FROM performance_schema.setup_instruments
    -> WHERE NAME LIKE '%wait/synch/mutex/innodb%';
+-----------------------------------------------------------+---------+-------+
| NAME                                                      | ENABLED | TIMED |
+-----------------------------------------------------------+---------+-------+
| wait/synch/mutex/innodb/commit_cond_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/innobase_share_mutex              | NO      | NO    |
| wait/synch/mutex/innodb/autoinc_mutex                     | NO      | NO    |
| wait/synch/mutex/innodb/autoinc_persisted_mutex           | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_flush_state_mutex        | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_LRU_list_mutex           | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_free_list_mutex          | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_zip_free_mutex           | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_zip_hash_mutex           | NO      | NO    |
| wait/synch/mutex/innodb/buf_pool_zip_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/cache_last_read_mutex             | NO      | NO    |
| wait/synch/mutex/innodb/dict_foreign_err_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/dict_persist_dirty_tables_mutex   | NO      | NO    |
| wait/synch/mutex/innodb/dict_sys_mutex                    | NO      | NO    |
| wait/synch/mutex/innodb/recalc_pool_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/fil_system_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/flush_list_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/fts_bg_threads_mutex              | NO      | NO    |
| wait/synch/mutex/innodb/fts_delete_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/fts_optimize_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/fts_doc_id_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/log_flush_order_mutex             | NO      | NO    |
| wait/synch/mutex/innodb/hash_table_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/ibuf_bitmap_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/ibuf_mutex                        | NO      | NO    |
| wait/synch/mutex/innodb/ibuf_pessimistic_insert_mutex     | NO      | NO    |
| wait/synch/mutex/innodb/log_sys_mutex                     | NO      | NO    |
| wait/synch/mutex/innodb/log_sys_write_mutex               | NO      | NO    |
| wait/synch/mutex/innodb/mutex_list_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/page_zip_stat_per_index_mutex     | NO      | NO    |
| wait/synch/mutex/innodb/purge_sys_pq_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/recv_sys_mutex                    | NO      | NO    |
| wait/synch/mutex/innodb/recv_writer_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/redo_rseg_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/noredo_rseg_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/rw_lock_list_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/rw_lock_mutex                     | NO      | NO    |
| wait/synch/mutex/innodb/srv_dict_tmpfile_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/srv_innodb_monitor_mutex          | NO      | NO    |
| wait/synch/mutex/innodb/srv_misc_tmpfile_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/srv_monitor_file_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/buf_dblwr_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/trx_undo_mutex                    | NO      | NO    |
| wait/synch/mutex/innodb/trx_pool_mutex                    | NO      | NO    |
| wait/synch/mutex/innodb/trx_pool_manager_mutex            | NO      | NO    |
| wait/synch/mutex/innodb/srv_sys_mutex                     | NO      | NO    |
| wait/synch/mutex/innodb/lock_mutex                        | NO      | NO    |
| wait/synch/mutex/innodb/lock_wait_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/trx_mutex                         | NO      | NO    |
| wait/synch/mutex/innodb/srv_threads_mutex                 | NO      | NO    |
```

```
| wait/synch/mutex/innodb/rtr_active_mutex                | NO      | NO    |
| wait/synch/mutex/innodb/rtr_match_mutex                 | NO      | NO    |
| wait/synch/mutex/innodb/rtr_path_mutex                  | NO      | NO    |
| wait/synch/mutex/innodb/rtr_ssn_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/trx_sys_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/zip_pad_mutex                   | NO      | NO    |
| wait/synch/mutex/innodb/master_key_id_mutex             | NO      | NO    |
+---------------------------------------------------------+---------+-------+
```

2. Some `InnoDB` mutex instances are created at server startup and are only instrumented if the associated instrument is also enabled at server startup. To ensure that all `InnoDB` mutex instances are instrumented and enabled, add the following `performance-schema-instrument` rule to your MySQL configuration file:

```
performance-schema-instrument='wait/synch/mutex/innodb/%=ON'
```

If you do not require wait event data for all `InnoDB` mutexes, you can disable specific instruments by adding additional `performance-schema-instrument` rules to your MySQL configuration file. For example, to disable `InnoDB` mutex wait event instruments related to full-text search, add the following rule:

```
performance-schema-instrument='wait/synch/mutex/innodb/fts%=OFF'
```

Note

Rules with a longer prefix such as `wait/synch/mutex/innodb/fts%` take precedence over rules with shorter prefixes such as `wait/synch/mutex/innodb/%`.

After adding the `performance-schema-instrument` rules to your configuration file, restart the server. All the `InnoDB` mutexes except for those related to full text search are enabled. To verify, query the `setup_instruments` table. The `ENABLED` and `TIMED` columns should be set to `YES` for the instruments that you enabled.

```
mysql> SELECT * FROM performance_schema.setup_instruments
    -> WHERE NAME LIKE '%wait/synch/mutex/innodb%';
+---------------------------------------------------------+---------+-------+
| NAME                                                    | ENABLED | TIMED |
+---------------------------------------------------------+---------+-------+
| wait/synch/mutex/innodb/commit_cond_mutex               | YES     | YES   |
| wait/synch/mutex/innodb/innobase_share_mutex            | YES     | YES   |
| wait/synch/mutex/innodb/autoinc_mutex                   | YES     | YES   |
...
| wait/synch/mutex/innodb/master_key_id_mutex             | YES     | YES   |
+---------------------------------------------------------+---------+-------+
49 rows in set (0.00 sec)
```

3. Enable wait event consumers by updating the `setup_consumers` table. Wait event consumers are disabled by default.

```
mysql> UPDATE performance_schema.setup_consumers SET enabled = 'YES'
    -> WHERE name like 'events_waits%';
Query OK, 3 rows affected (0.00 sec)
Rows matched: 3  Changed: 3  Warnings: 0
```

You can verify that wait event consumers are enabled by querying the `setup_consumers` table. The `events_waits_current`, `events_waits_history`, and `events_waits_history_long` consumers should be enabled.

```
mysql> SELECT * FROM performance_schema.setup_consumers;
+--------------------------------+---------+
| NAME                           | ENABLED |
+--------------------------------+---------+
| events_stages_current          | NO      |
| events_stages_history          | NO      |
| events_stages_history_long     | NO      |
```

```
| events_statements_current        | YES     |
| events_statements_history        | YES     |
| events_statements_history_long   | NO      |
| events_transactions_current      | YES     |
| events_transactions_history      | YES     |
| events_transactions_history_long | NO      |
| events_waits_current             | YES     |
| events_waits_history             | YES     |
| events_waits_history_long        | YES     |
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| statements_digest                | YES     |
+----------------------------------+---------+
15 rows in set (0.00 sec)
```

4. Once instruments and consumers are enabled, run the workload that you want to monitor. In this example, the `mysqlslap` load emulation client is used to simulate a workload.

```
shell> ./mysqlslap --auto-generate-sql --concurrency=100 --iterations=10
    -> --number-of-queries=1000 --number-char-cols=6 --number-int-cols=6;
```

5. Query the wait event data. In this example, wait event data is queried from the `events_waits_summary_global_by_event_name` table which aggregates data found in the `events_waits_current`, `events_waits_history`, and `events_waits_history_long` tables. Data is summarized by event name (`EVENT_NAME`), which is the name of the instrument that produced the event. Summarized data includes:

   - `COUNT_STAR`

     The number of summarized wait events.

   - `SUM_TIMER_WAIT`

     The total wait time of the summarized timed wait events.

   - `MIN_TIMER_WAIT`

     The minimum wait time of the summarized timed wait events.

   - `AVG_TIMER_WAIT`

     The average wait time of the summarized timed wait events.

   - `MAX_TIMER_WAIT`

     The maximum wait time of the summarized timed wait events.

   The following query returns the instrument name (`EVENT_NAME`), the number of wait events (`COUNT_STAR`), and the total wait time for the events for that instrument (`SUM_TIMER_WAIT`). Because waits are timed in picoseconds (trillionths of a second) by default, wait times are divided by 1000000000 to show wait times in milliseconds. Data is presented in descending order, by the number of summarized wait events (`COUNT_STAR`). You can adjust the `ORDER BY` clause to order the data by total wait time.

```
mysql> SELECT EVENT_NAME, COUNT_STAR, SUM_TIMER_WAIT/1000000000 SUM_TIMER_WAIT_MS
    -> FROM performance_schema.events_waits_summary_global_by_event_name
    -> WHERE SUM_TIMER_WAIT > 0 AND EVENT_NAME LIKE 'wait/synch/mutex/innodb/%'
    -> ORDER BY COUNT_STAR DESC;
+-----------------------------------------------------+------------+-------------------+
| EVENT_NAME                                          | COUNT_STAR | SUM_TIMER_WAIT_MS |
+-----------------------------------------------------+------------+-------------------+
| wait/synch/mutex/innodb/trx_mutex                   |     201111 |           23.4719 |
| wait/synch/mutex/innodb/fil_system_mutex            |      62244 |            9.6426 |
| wait/synch/mutex/innodb/redo_rseg_mutex             |      48238 |            3.1135 |
| wait/synch/mutex/innodb/log_sys_mutex               |      46113 |            2.0434 |
```

```
| wait/synch/mutex/innodb/trx_sys_mutex                    |      35134 |          1068.1588 |
| wait/synch/mutex/innodb/lock_mutex                       |      34872 |          1039.2589 |
| wait/synch/mutex/innodb/log_sys_write_mutex              |      17805 |          1526.0490 |
| wait/synch/mutex/innodb/dict_sys_mutex                   |      14912 |          1606.7348 |
| wait/synch/mutex/innodb/trx_undo_mutex                   |      10634 |             1.1424 |
| wait/synch/mutex/innodb/rw_lock_list_mutex               |       8538 |             0.1960 |
| wait/synch/mutex/innodb/buf_pool_free_list_mutex         |       5961 |             0.6473 |
| wait/synch/mutex/innodb/trx_pool_mutex                   |       4885 |          8821.7496 |
| wait/synch/mutex/innodb/buf_pool_LRU_list_mutex          |       4364 |             0.2077 |
| wait/synch/mutex/innodb/innobase_share_mutex             |       3212 |             0.2650 |
| wait/synch/mutex/innodb/flush_list_mutex                 |       3178 |             0.2349 |
| wait/synch/mutex/innodb/trx_pool_manager_mutex           |       2495 |             0.1310 |
| wait/synch/mutex/innodb/buf_pool_flush_state_mutex       |       1318 |             0.2161 |
| wait/synch/mutex/innodb/log_flush_order_mutex            |       1250 |             0.0893 |
| wait/synch/mutex/innodb/buf_dblwr_mutex                  |        951 |             0.0918 |
| wait/synch/mutex/innodb/recalc_pool_mutex                |        670 |             0.0942 |
| wait/synch/mutex/innodb/dict_persist_dirty_tables_mutex  |        345 |             0.0414 |
| wait/synch/mutex/innodb/lock_wait_mutex                  |        303 |             0.1565 |
| wait/synch/mutex/innodb/autoinc_mutex                    |        196 |             0.0213 |
| wait/synch/mutex/innodb/autoinc_persisted_mutex          |        196 |             0.0175 |
| wait/synch/mutex/innodb/purge_sys_pq_mutex               |        117 |             0.0308 |
| wait/synch/mutex/innodb/srv_sys_mutex                    |         94 |             0.0077 |
| wait/synch/mutex/innodb/ibuf_mutex                       |         22 |             0.0086 |
| wait/synch/mutex/innodb/recv_sys_mutex                   |         12 |             0.0008 |
| wait/synch/mutex/innodb/srv_innodb_monitor_mutex         |          4 |             0.0009 |
| wait/synch/mutex/innodb/recv_writer_mutex                |          1 |             0.0005 |
+----------------------------------------------------------+------------+--------------------+
```

Note

The preceding result set includes wait event data produced during the startup process. To exclude this data, you can truncate the `events_waits_summary_global_by_event_name` table immediately after startup and before running your workload. However, the truncate operation itself may produce a negligible amount wait event data.

```
mysql> TRUNCATE performance_schema.events_waits_summary_global_by_event_name;
```

15.16 InnoDB Monitors

`InnoDB` monitors provide information about the `InnoDB` internal state. This information is useful for performance tuning.

15.16.1 InnoDB Monitor Types

There are two types of `InnoDB` monitor:

- The standard `InnoDB` Monitor displays the following types of information:
  - Work done by the main background thread
  - Semaphore waits
  - Data about the most recent foreign key and deadlock errors
  - Lock waits for transactions
  - Table and record locks held by active transactions
  - Pending I/O operations and related statistics
  - Insert buffer and adaptive hash index statistics
  - Redo log data
  - Buffer pool statistics

  - Row operation data

- The InnoDB Lock Monitor prints additional lock information as part of the standard InnoDB Monitor output.

15.16.2 Enabling InnoDB Monitors

When you enable InnoDB monitors for periodic output, InnoDB writes the output to the mysqld server standard error output (stderr). In this case, no output is sent to clients. When switched on, InnoDB monitors print data about every 15 seconds. Server output usually is directed to the error log (see Section 5.4.2, “The Error Log”). This data is useful in performance tuning. On Windows, start the server from a command prompt in a console window with the --console option if you want to direct the output to the window rather than to the error log.

InnoDB sends diagnostic output to stderr or to files rather than to stdout or fixed-size memory buffers, to avoid potential buffer overflows. As a side effect, the output of SHOW ENGINE INNODB STATUS is written to a status file in the MySQL data directory every fifteen seconds. The name of the file is innodb_status.*pid*, where *pid* is the server process ID. InnoDB removes the file when the server is shut down normally. If abnormal shutdowns have occurred, instances of these status files may be present and must be removed manually. Before removing them, examine them to see if they contain useful information about the cause of abnormal shutdowns. The innodb_status.*pid* file is only created if the innodb-status-file configuration option is enabled. It is disabled by default.

InnoDB monitors should only be enabled when you actually want to see monitor information because output generation causes some performance decrement. Also, if you enable monitor output, your error log may become quite large if you forget to disable it later.

Note

To assist with troubleshooting, InnoDB temporarily enables standard InnoDB Monitor output under certain conditions. For more information, see Section 15.20, “InnoDB Troubleshooting”.

InnoDB monitor output begins with a header containing a timestamp and the monitor name. For example:

```
===================================
2014-10-16 18:37:29 0x7fc2a95c1700 INNODB MONITOR OUTPUT
===================================
```

The header for the standard InnoDB Monitor (INNODB MONITOR OUTPUT) is also used for the Lock Monitor because the latter produces the same output with the addition of extra lock information.

The innodb_status_output and innodb_status_output_locks system variables are used to enable the standard InnoDB Monitor and InnoDB Lock Monitor.

The PROCESS privilege is required to enable and disable InnoDB Monitors.

Enabling the Standard InnoDB Monitor

Enable the standard InnoDB Monitor by setting the innodb_status_output system variable to ON.

```
set GLOBAL innodb_status_output=ON;
```

To disable the standard InnoDB Monitor, set innodb_status_output to OFF.

When you shut down the server, the innodb_status_output variable is set to the default OFF value.

Obtaining Standard InnoDB Monitor Output On Demand

As an alternative to enabling the standard InnoDB Monitor for periodic output, you can obtain standard InnoDB Monitor output on demand using the SHOW ENGINE INNODB STATUS SQL statement, which fetches the output to your client program. If you are using the mysql interactive client, the output is more readable if you replace the usual semicolon statement terminator with \G:

```
mysql> SHOW ENGINE INNODB STATUS\G
```

SHOW ENGINE INNODB STATUS output also includes InnoDB Lock Monitor data if the InnoDB Lock Monitor is enabled.

Enabling the InnoDB Lock Monitor

InnoDB Lock Monitor data is printed with the InnoDB Standard Monitor output. Both the InnoDB Standard Monitor and InnoDB Lock Monitor must be enabled to have InnoDB Lock Monitor data printed periodically.

To enable the InnoDB Lock Monitor, set the innodb_status_output_locks system variable to ON. Both the InnoDB standard Monitor and InnoDB Lock Monitor must be enabled to have InnoDB Lock Monitor data printed periodically:

```
set GLOBAL innodb_status_output=ON;
set GLOBAL innodb_status_output_locks=ON;
```

To disable the InnoDB Lock Monitor, set innodb_status_output_locks to OFF. Set innodb_status_output to OFF to also disable the InnoDB Standard Monitor.

When you shut down the server, the innodb_status_output and innodb_status_output_locks variables are set to the default OFF value.

Note

To enable the InnoDB Lock Monitor for SHOW ENGINE INNODB STATUS output, you are only required to enable innodb_status_output_locks.

15.16.3 InnoDB Standard Monitor and Lock Monitor Output

The Lock Monitor is the same as the Standard Monitor except that it includes additional lock information. Enabling either monitor for periodic output turns on the same output stream, but the stream includes extra information if the Lock Monitor is enabled. For example, if you enable the Standard Monitor and Lock Monitor, that turns on a single output stream. The stream includes extra lock information until you disable the Lock Monitor.

Standard Monitor output is limited to 1MB when produced using the SHOW ENGINE INNODB STATUS statement. This limit does not apply to output written to server standard error output (stderr).

Example Standard Monitor output:

```
mysql> SHOW ENGINE INNODB STATUS\G
*************************** 1. row ***************************
  Type: InnoDB
  Name:
Status:
=====================================
2014-10-16 18:37:29 0x7fc2a95c1700 INNODB MONITOR OUTPUT
=====================================
Per second averages calculated from the last 20 seconds
-----------------
BACKGROUND THREAD
-----------------
srv_master_thread loops: 38 srv_active, 0 srv_shutdown, 252 srv_idle
srv_master_thread log flush and writes: 290
```

```
----------
SEMAPHORES
----------
OS WAIT ARRAY INFO: reservation count 119
OS WAIT ARRAY INFO: signal count 103
Mutex spin waits 0, rounds 0, OS waits 0
RW-shared spins 38, rounds 76, OS waits 38
RW-excl spins 2, rounds 9383715, OS waits 3
RW-sx spins 0, rounds 0, OS waits 0
Spin rounds per wait: 0.00 mutex, 2.00 RW-shared, 4691857.50 RW-excl,
0.00 RW-sx
------------------------
LATEST FOREIGN KEY ERROR
------------------------
2014-10-16 18:35:18 0x7fc2a95c1700 Transaction:
TRANSACTION 1814, ACTIVE 0 sec inserting
mysql tables in use 1, locked 1
4 lock struct(s), heap size 1136, 3 row lock(s), undo log entries 3
MySQL thread id 2, OS thread handle 140474041767680, query id 74 localhost
root update
INSERT INTO child VALUES
    (NULL, 1)
    , (NULL, 2)
    , (NULL, 3)
    , (NULL, 4)
    , (NULL, 5)
    , (NULL, 6)
Foreign key constraint fails for table `mysql`.`child`:
,
  CONSTRAINT `child_ibfk_1` FOREIGN KEY (`parent_id`) REFERENCES `parent`
  (`id`) ON DELETE CASCADE ON UPDATE CASCADE
Trying to add in child table, in index par_ind tuple:
DATA TUPLE: 2 fields;
 0: len 4; hex 80000003; asc     ;;
 1: len 4; hex 80000003; asc     ;;

But in parent table `mysql`.`parent`, in index PRIMARY,
the closest match we can find is record:
PHYSICAL RECORD: n_fields 3; compact format; info bits 0
 0: len 4; hex 80000004; asc     ;;
 1: len 6; hex 00000000070a; asc       ;;
 2: len 7; hex aa0000011d0134; asc       4;;

------------------------
LATEST DETECTED DEADLOCK
------------------------
2014-10-16 18:36:30 0x7fc2a95c1700
*** (1) TRANSACTION:
TRANSACTION 1824, ACTIVE 9 sec starting index read
mysql tables in use 1, locked 1
LOCK WAIT 2 lock struct(s), heap size 1136, 1 row lock(s)
MySQL thread id 3, OS thread handle 140474041501440, query id 80 localhost
root updating
DELETE FROM t WHERE i = 1
*** (1) WAITING FOR THIS LOCK TO BE GRANTED:
RECORD LOCKS space id 35 page no 3 n bits 72 index GEN_CLUST_INDEX of table
`mysql`.`t` trx id 1824 lock_mode X waiting
Record lock, heap no 2 PHYSICAL RECORD: n_fields 4; compact format; info
bits 0
 0: len 6; hex 000000000200; asc       ;;
 1: len 6; hex 00000000071f; asc       ;;
 2: len 7; hex b80000012b0110; asc     +  ;;
 3: len 4; hex 80000001; asc     ;;

*** (2) TRANSACTION:
TRANSACTION 1825, ACTIVE 29 sec starting index read
mysql tables in use 1, locked 1
4 lock struct(s), heap size 1136, 3 row lock(s)
MySQL thread id 2, OS thread handle 140474041767680, query id 81 localhost
root updating
DELETE FROM t WHERE i = 1
*** (2) HOLDS THE LOCK(S):
```

```
RECORD LOCKS space id 35 page no 3 n bits 72 index GEN_CLUST_INDEX of table
`mysql`.`t` trx id 1825 lock mode S
Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info
bits 0
 0: len 8; hex 73757072656d756d; asc supremum;;

Record lock, heap no 2 PHYSICAL RECORD: n_fields 4; compact format; info bits 0
 0: len 6; hex 000000000200; asc       ;;
 1: len 6; hex 00000000071f; asc       ;;
 2: len 7; hex b80000012b0110; asc     +  ;;
 3: len 4; hex 80000001; asc     ;;

*** (2) WAITING FOR THIS LOCK TO BE GRANTED:
RECORD LOCKS space id 35 page no 3 n bits 72 index GEN_CLUST_INDEX of table
`mysql`.`t` trx id 1825 lock_mode X waiting
Record lock, heap no 2 PHYSICAL RECORD: n_fields 4; compact format; info
bits 0
 0: len 6; hex 000000000200; asc       ;;
 1: len 6; hex 00000000071f; asc       ;;
 2: len 7; hex b80000012b0110; asc     +  ;;
 3: len 4; hex 80000001; asc     ;;

*** WE ROLL BACK TRANSACTION (1)
------------
TRANSACTIONS
------------
Trx id counter 1950
Purge done for trx's n:o < 1933 undo n:o < 0 state: running but idle
History list length 23
LIST OF TRANSACTIONS FOR EACH SESSION:
---TRANSACTION 421949033065200, not started
0 lock struct(s), heap size 1136, 0 row lock(s)
---TRANSACTION 421949033064280, not started
0 lock struct(s), heap size 1136, 0 row lock(s)
---TRANSACTION 1949, ACTIVE 0 sec inserting
mysql tables in use 1, locked 1
8 lock struct(s), heap size 1136, 1850 row lock(s), undo log entries 17415
MySQL thread id 4, OS thread handle 140474041235200, query id 176 localhost
root update
INSERT INTO `salaries` VALUES (55723,39746,'1997-02-25','1998-02-25'),
(55723,40758,'1998-02-25','1999-02-25'),(55723,44559,'1999-02-25','2000-02-25'),
(55723,44081,'2000-02-25','2001-02-24'),(55723,44112,'2001-02-24','2001-08-16'),
(55724,46461,'1996-12-06','1997-12-06'),(55724,48916,'1997-12-06','1998-12-06'),
(55724,51269,'1998-12-06','1999-12-06'),(55724,51932,'1999-12-06','2000-12-05'),
(55724,52617,'2000-12-05','2001-12-05'),(55724,56658,'2001-12-05','9999-01-01'),
(55725,40000,'1993-01-30','1994-01-30'),(55725,41472,'1994-01-30','1995-01-30'),
(55725,45293,'1995-01-30','1996-01-30'),(55725,473
--------
FILE I/O
--------
I/O thread 0 state: waiting for completed aio requests (insert buffer thread)
I/O thread 1 state: waiting for completed aio requests (log thread)
I/O thread 2 state: waiting for completed aio requests (read thread)
I/O thread 3 state: waiting for completed aio requests (read thread)
I/O thread 4 state: waiting for completed aio requests (read thread)
I/O thread 5 state: waiting for completed aio requests (read thread)
I/O thread 6 state: waiting for completed aio requests (write thread)
I/O thread 7 state: waiting for completed aio requests (write thread)
I/O thread 8 state: waiting for completed aio requests (write thread)
I/O thread 9 state: waiting for completed aio requests (write thread)
Pending normal aio reads: 0 [0, 0, 0, 0] , aio writes: 0 [0, 0, 0, 0] ,
 ibuf aio reads: 0, log i/o's: 0, sync i/o's: 0
Pending flushes (fsync) log: 0; buffer pool: 0
224 OS file reads, 5770 OS file writes, 803 OS fsyncs
0.00 reads/s, 0 avg bytes/read, 264.84 writes/s, 23.05 fsyncs/s
-------------------------------------
INSERT BUFFER AND ADAPTIVE HASH INDEX
-------------------------------------
Ibuf: size 1, free list len 0, seg size 2, 0 merges
merged operations:
 insert 0, delete mark 0, delete 0
discarded operations:
```

```
 insert 0, delete mark 0, delete 0
Hash table size 4425293, node heap has 444 buffer(s)
68015.25 hash searches/s, 106259.24 non-hash searches/s
---
LOG
---
Log sequence number 165913808
Log flushed up to   164814979
Pages flushed up to 141544038
Last checkpoint at  130503656
0 pending log flushes, 0 pending chkp writes
258 log i/o's done, 6.65 log i/o's/second
----------------------
BUFFER POOL AND MEMORY
----------------------
Total large memory allocated 2198863872
Dictionary memory allocated 776332
Buffer pool size   131072
Free buffers       124908
Database pages     5720
Old database pages 2071
Modified db pages  910
Pending reads 0
Pending writes: LRU 0, flush list 0, single page 0
Pages made young 4, not young 0
0.10 youngs/s, 0.00 non-youngs/s
Pages read 197, created 5523, written 5060
0.00 reads/s, 190.89 creates/s, 244.94 writes/s
Buffer pool hit rate 1000 / 1000, young-making rate 0 / 1000 not
0 / 1000
Pages read ahead 0.00/s, evicted without access 0.00/s, Random read
ahead 0.00/s
LRU len: 5720, unzip_LRU len: 0
I/O sum[0]:cur[0], unzip sum[0]:cur[0]
----------------------
INDIVIDUAL BUFFER POOL INFO
----------------------
---BUFFER POOL 0
Buffer pool size   65536
Free buffers       62412
Database pages     2899
Old database pages 1050
Modified db pages  449
Pending reads 0
Pending writes: LRU 0, flush list 0, single page 0
Pages made young 3, not young 0
0.05 youngs/s, 0.00 non-youngs/s
Pages read 107, created 2792, written 2586
0.00 reads/s, 92.65 creates/s, 122.89 writes/s
Buffer pool hit rate 1000 / 1000, young-making rate 0 / 1000 not 0 / 1000
Pages read ahead 0.00/s, evicted without access 0.00/s, Random read ahead
0.00/s
LRU len: 2899, unzip_LRU len: 0
I/O sum[0]:cur[0], unzip sum[0]:cur[0]
---BUFFER POOL 1
Buffer pool size   65536
Free buffers       62496
Database pages     2821
Old database pages 1021
Modified db pages  461
Pending reads 0
Pending writes: LRU 0, flush list 0, single page 0
Pages made young 1, not young 0
0.05 youngs/s, 0.00 non-youngs/s
Pages read 90, created 2731, written 2474
0.00 reads/s, 98.25 creates/s, 122.04 writes/s
Buffer pool hit rate 1000 / 1000, young-making rate 0 / 1000 not 0 / 1000
Pages read ahead 0.00/s, evicted without access 0.00/s, Random read ahead
0.00/s
LRU len: 2821, unzip_LRU len: 0
I/O sum[0]:cur[0], unzip sum[0]:cur[0]
--------------
```

```
ROW OPERATIONS
-------------
0 queries inside InnoDB, 0 queries in queue
0 read views open inside InnoDB
Process ID=35909, Main thread ID=140471692396288, state: sleeping
Number of rows inserted 1526363, updated 0, deleted 3, read 11
52671.72 inserts/s, 0.00 updates/s, 0.00 deletes/s, 0.00 reads/s
----------------------------
END OF INNODB MONITOR OUTPUT
============================
```

Standard Monitor Output Sections

For a description of each metric reported by the Standard Monitor, refer to the Metrics chapter in the Oracle Enterprise Manager for MySQL Database User's Guide.

- `Status`

  This section shows the timestamp, the monitor name, and the number of seconds that per-second averages are based on. The number of seconds is the elapsed time between the current time and the last time `InnoDB` Monitor output was printed.

- `BACKGROUND THREAD`

  The `srv_master_thread` lines shows work done by the main background thread.

- `SEMAPHORES`

  This section reports threads waiting for a semaphore and statistics on how many times threads have needed a spin or a wait on a mutex or a rw-lock semaphore. A large number of threads waiting for semaphores may be a result of disk I/O, or contention problems inside `InnoDB`. Contention can be due to heavy parallelism of queries or problems in operating system thread scheduling. Setting the `innodb_thread_concurrency` system variable smaller than the default value might help in such situations. The `Spin rounds per wait` line shows the number of spinlock rounds per OS wait for a mutex.

  Mutex metrics are reported by `SHOW ENGINE INNODB MUTEX`.

- `LATEST FOREIGN KEY ERROR`

  This section provides information about the most recent foreign key constraint error. It is not present if no such error has occurred. The contents include the statement that failed as well as information about the constraint that failed and the referenced and referencing tables.

- `LATEST DETECTED DEADLOCK`

  This section provides information about the most recent deadlock. It is not present if no deadlock has occurred. The contents show which transactions are involved, the statement each was attempting to execute, the locks they have and need, and which transaction `InnoDB` decided to roll back to break the deadlock. The lock modes reported in this section are explained in Section 15.5.1, “InnoDB Locking”.

- `TRANSACTIONS`

  If this section reports lock waits, your applications might have lock contention. The output can also help to trace the reasons for transaction deadlocks.

- `FILE I/O`

  This section provides information about threads that `InnoDB` uses to perform various types of I/O. The first few of these are dedicated to general `InnoDB` processing. The contents also display information for pending I/O operations and statistics for I/O performance.

The number of these threads are controlled by the `innodb_read_io_threads` and `innodb_write_io_threads` parameters. See Section 15.13, “InnoDB Startup Options and System Variables”.

- `INSERT BUFFER AND ADAPTIVE HASH INDEX`

  This section shows the status of the `InnoDB` insert buffer (also referred to as the change buffer) and the adaptive hash index.

  For related information, see Section 15.4.2, “Change Buffer”, and Section 15.4.3, “Adaptive Hash Index”.

- `LOG`

  This section displays information about the `InnoDB` log. The contents include the current log sequence number, how far the log has been flushed to disk, and the position at which `InnoDB` last took a checkpoint. (See Section 15.11.3, “InnoDB Checkpoints”.) The section also displays information about pending writes and write performance statistics.

- `BUFFER POOL AND MEMORY`

  This section gives you statistics on pages read and written. You can calculate from these numbers how many data file I/O operations your queries currently are doing.

  For buffer pool statistics descriptions, see Section 15.6.3.9, “Monitoring the Buffer Pool Using the InnoDB Standard Monitor”. For additional information about the operation of the buffer pool, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

- `ROW OPERATIONS`

  This section shows what the main thread is doing, including the number and performance rate for each type of row operation.

15.17 InnoDB Backup and Recovery

This section covers topics related to `InnoDB` backup and recovery.

- For information about backup techniques applicable to `InnoDB`, see Section 15.17.1, “InnoDB Backup”.
- For information about point-in-time recovery, recovery from disk failure or corruption, and how `InnoDB` performs crash recovery, see Section 15.17.2, “InnoDB Recovery”.

15.17.1 InnoDB Backup

The key to safe database management is making regular backups. Depending on your data volume, number of MySQL servers, and database workload, you can use these backup techniques, alone or in combination: hot backup with MySQL Enterprise Backup; cold backup by copying files while the MySQL server is shut down; logical backup with `mysqldump` for smaller data volumes or to record the structure of schema objects. Hot and cold backups are physical backups that copy actual data files, which can be used directly by the `mysqld` server for faster restore.

Hot Backups

The `mysqlbackup` command, part of the MySQL Enterprise Backup component, lets you back up a running MySQL instance, including `InnoDB` tables, with minimal disruption to operations while producing a consistent snapshot of the database. When `mysqlbackup` is copying `InnoDB` tables, reads and writes to `InnoDB` tables can continue. MySQL Enterprise Backup can also create compressed backup files, and back up subsets of tables and databases. In conjunction with the MySQL

binary log, users can perform point-in-time recovery. MySQL Enterprise Backup is part of the MySQL Enterprise subscription. For more details, see Section 29.2, “MySQL Enterprise Backup Overview”.

Cold Backups

If you can shut down the MySQL server, you can make a physical backup that consists of all files used by `InnoDB` to manage its tables. Use the following procedure:

1. Perform a slow shutdown of the MySQL server and make sure that it stops without errors.
2. Copy all `InnoDB` data files (`ibdata` files and `.ibd` files) into a safe place.
3. Copy all `InnoDB` log files (`ib_logfile` files) to a safe place.
4. Copy your `my.cnf` configuration file or files to a safe place.

Logical Backups Using mysqldump

In addition to physical backups, it is recommended that you regularly create logical backups by dumping your tables using `mysqldump`. A binary file might be corrupted without you noticing it. Dumped tables are stored into text files that are human-readable, so spotting table corruption becomes easier. Also, because the format is simpler, the chance for serious data corruption is smaller. `mysqldump` also has a `--single-transaction` option for making a consistent snapshot without locking out other clients. See Section 7.3.1, “Establishing a Backup Policy”.

Replication works with `InnoDB` tables, so you can use MySQL replication capabilities to keep a copy of your database at database sites requiring high availability. See Section 15.18, “InnoDB and MySQL Replication”.

15.17.2 InnoDB Recovery

This section describes `InnoDB` recovery. Topics include:

- Point-in-Time Recovery
- Recovery from Data Corruption or Disk Failure
- InnoDB Crash Recovery
- Tablespace Discovery During Crash Recovery
- Lost or Corrupted Tablespace Map Files

Point-in-Time Recovery

To recover an `InnoDB` database to the present from the time at which the physical backup was made, you must run MySQL server with binary logging enabled, even before taking the backup. To achieve point-in-time recovery after restoring a backup, you can apply changes from the binary log that occurred after the backup was made. See Section 7.5, “Point-in-Time (Incremental) Recovery Using the Binary Log”.

Recovery from Data Corruption or Disk Failure

If your database becomes corrupted or disk failure occurs, you must perform the recovery using a backup. In the case of corruption, first find a backup that is not corrupted. After restoring the base backup, do a point-in-time recovery from the binary log files using `mysqlbinlog` and `mysql` to restore the changes that occurred after the backup was made.

In some cases of database corruption, it is enough to dump, drop, and re-create one or a few corrupt tables. You can use the `CHECK TABLE` statement to check whether a table is corrupt, although `CHECK TABLE` naturally cannot detect every possible kind of corruption.

In some cases, apparent database page corruption is actually due to the operating system corrupting its own file cache, and the data on disk may be okay. It is best to try restarting the computer first. Doing so may eliminate errors that appeared to be database page corruption. If MySQL still has trouble starting because of `InnoDB` consistency problems, see Section 15.20.2, “Forcing InnoDB Recovery” for steps to start the instance in recovery mode, which permits you to dump the data.

InnoDB Crash Recovery

To recover from a MySQL server crash, the only requirement is to restart the MySQL server. `InnoDB` automatically checks the logs and performs a roll-forward of the database to the present. `InnoDB` automatically rolls back uncommitted transactions that were present at the time of the crash. During recovery, `mysqld` displays output similar to this:

```
InnoDB: The log sequence number 664050266 in the system tablespace does not match
the log sequence number 685111586 in the ib_logfiles!
InnoDB: Database was not shutdown normally!
InnoDB: Starting crash recovery.
InnoDB: Using 'tablespaces.open.2' max LSN: 664075228
InnoDB: Doing recovery: scanned up to log sequence number 690354176
InnoDB: Doing recovery: scanned up to log sequence number 695597056
InnoDB: Doing recovery: scanned up to log sequence number 700839936
InnoDB: Doing recovery: scanned up to log sequence number 706082816
InnoDB: Doing recovery: scanned up to log sequence number 711325696
InnoDB: Doing recovery: scanned up to log sequence number 713458156
InnoDB: Applying a batch of 1467 redo log records ...
InnoDB: 10%
InnoDB: 20%
InnoDB: 30%
InnoDB: 40%
InnoDB: 50%
InnoDB: 60%
InnoDB: 70%
InnoDB: 80%
InnoDB: 90%
InnoDB: 100%
InnoDB: Apply batch completed!
InnoDB: 1 transaction(s) which must be rolled back or cleaned up in total 561887 row
operations to undo
InnoDB: Trx id counter is 4096
...
InnoDB: 8.0.1 started; log sequence number 713458156
InnoDB: Waiting for purge to start
InnoDB: Starting in background the rollback of uncommitted transactions
InnoDB: Rolling back trx with id 3596, 561887 rows to undo
...
./mysqld: ready for connections....
```

`InnoDB` crash recovery consists of several steps:

- Tablespace discovery

  Tablespace discovery is the process that `InnoDB` uses to identify tablespaces that require redo log application. See Tablespace Discovery During Crash Recovery.

- Redo log application

  Redo log application is performed during initialization, before accepting any connections. If all changes are flushed from the buffer pool to the tablespaces (`ibdata*` and `*.ibd` files) at the time of the shutdown or crash, redo log application is skipped. `InnoDB` also skips redo log application if redo log files are missing at startup.

  - The current maximum auto-increment counter value is written to the redo log each time the value changes, which makes it crash-safe. During recovery, `InnoDB` scans the redo log to collect counter value changes and applies the changes to the in-memory table object.

For more information about how InnoDB handles auto-increment values, see Section 15.8.1.5, “AUTO_INCREMENT Handling in InnoDB”, and InnoDB AUTO_INCREMENT Counter Initialization.

- When encountering index tree corruption, InnoDB writes a corruption flag to the redo log, which makes the corruption flag crash-safe. InnoDB also writes in-memory corruption flag data to an engine-private system table on each checkpoint. During recovery, InnoDB reads corruption flags from both locations and merges results before marking in-memory table and index objects as corrupt.
- Removing redo logs to speed up recovery is not recommended, even if some data loss is acceptable. Removing redo logs should only be considered after a clean shutdown, with innodb_fast_shutdown set to 0 or 1.

- Roll back of incomplete transactions

  Incomplete transactions are any transactions that were active at the time of crash or fast shutdown. The time it takes to roll back an incomplete transaction can be three or four times the amount of time a transaction is active before it is interrupted, depending on server load.

  You cannot cancel transactions that are being rolled back. In extreme cases, when rolling back transactions is expected to take an exceptionally long time, it may be faster to start InnoDB with an innodb_force_recovery setting of 3 or greater. See Section 15.20.2, “Forcing InnoDB Recovery”.

- Change buffer merge

  Applying changes from the change buffer (part of the system tablespace) to leaf pages of secondary indexes, as the index pages are read to the buffer pool.

- Purge

  Deleting delete-marked records that are no longer visible to active transactions.

The steps that follow redo log application do not depend on the redo log (other than for logging the writes) and are performed in parallel with normal processing. Of these, only rollback of incomplete transactions is special to crash recovery. The insert buffer merge and the purge are performed during normal processing.

After redo log application, InnoDB attempts to accept connections as early as possible, to reduce downtime. As part of crash recovery, InnoDB rolls back transactions that were not committed or in XA PREPARE state when the server crashed. The rollback is performed by a background thread, executed in parallel with transactions from new connections. Until the rollback operation is completed, new connections may encounter locking conflicts with recovered transactions.

In most situations, even if the MySQL server was killed unexpectedly in the middle of heavy activity, the recovery process happens automatically and no action is required of the DBA. If a hardware failure or severe system error corrupted InnoDB data, MySQL might refuse to start. In this case, see Section 15.20.2, “Forcing InnoDB Recovery”.

For information about the binary log and InnoDB crash recovery, see Section 5.4.4, “The Binary Log”.

Tablespace Discovery During Crash Recovery

If, during recovery, InnoDB encounters redo logs written since the last checkpoint, the redo logs must be applied to affected tablespaces. The process that identifies affected tablespaces during recovery is referred to as *tablespace discovery*.

In MySQL 8.0.1 and earlier, tablespace discovery is performed by scanning redo logs for MLOG_FILE_NAME records that contain tablespace ID and file name data. For information about this

method of tablespace discovery, refer to Tablespace Discovery During Crash Recovery in the MySQL 5.7 Reference Manual.

As of MySQL 8.0.2, tablespace discovery is performed using tablespace map files that map tablespace IDs to tablespace file names. Tablespace map files are stored in the `innodb_data_home_dir` directory. If `innodb_data_home_dir` is not configured, the default location is the MySQL data directory (`datadir`).

There are two tablespace map files (`tablespaces.open.1` and `tablespaces.open.2`) that are written to in circular fashion. Tablespace map files are only used during recovery. The files are ignored during normal startup.

In case of lost or corrupted tablespace map files, see Lost or Corrupted Tablespace Map Files.

Lost or Corrupted Tablespace Map Files

If tablespace map files, introduced in MySQL 8.0.2, are lost or corrupted, the `innodb_scan_directories` option can be used to specify tablespace file directories at startup. This option causes `InnoDB` to read the first page of each tablespace file in the specified directories and recreate tablespace map files so that the recovery process can apply redo logs.

This option can also be used to specify the directory path of a missing tablespace file. For example, if recovery reports an error due to a missing tablespace file, you can configure `innodb_scan_directories` to search for the tablespace file in a specific directory.

`innodb_scan_directories` may be specified as an option in a startup command or in a MySQL configuration file. Quotes are used around the argument value because otherwise a semicolon (;) is interpreted as a special character by some command interpreters. (Unix shells treat it as a command terminator, for example.)

Startup command:

```
mysqld --innodb-scan-directories="directory_path_1;directory_path_2"
```

MySQL configuration file:

```
[mysqld]
innodb_scan_directories="directory_path_1;directory_path_2"
```

When `innodb_scan_directories` is specified at startup, the `InnoDB` startup process prints messages similar to the following, reporting the directories that were scanned and the number of tablespace files found:

```
InnoDB: Directories to scan 'directory_path_1;directory_path_2'
InnoDB: Scanning 'directory_path_1'
InnoDB: Scanning 'directory_path_2'
InnoDB: Found 10 '.ibd' file(s)
```

15.18 InnoDB and MySQL Replication

MySQL replication works for `InnoDB` tables as it does for `MyISAM` tables. It is also possible to use replication in a way where the storage engine on the slave is not the same as the original storage engine on the master. For example, you can replicate modifications to an `InnoDB` table on the master to a `MyISAM` table on the slave. For more information see, Section 18.3.4, “Using Replication with Different Master and Slave Storage Engines”.

For information about setting up a new slave for a master, see Section 18.1.2.6, “Setting Up Replication Slaves”, and Section 18.1.2.5, “Choosing a Method for Data Snapshots”. To make a new slave without taking down the master or an existing slave, use the MySQL Enterprise Backup product.

Transactions that fail on the master do not affect replication at all. MySQL replication is based on the binary log where MySQL writes SQL statements that modify data. A transaction that fails (for example,

because of a foreign key violation, or because it is rolled back) is not written to the binary log, so it is not sent to slaves. See Section 13.3.1, “START TRANSACTION, COMMIT, and ROLLBACK Syntax”.

Replication and CASCADE. Cascading actions for `InnoDB` tables on the master are replicated on the slave *only* if the tables sharing the foreign key relation use `InnoDB` on both the master and slave. This is true whether you are using statement-based or row-based replication. Suppose that you have started replication, and then create two tables on the master using the following `CREATE TABLE` statements:

```
CREATE TABLE fc1 (
    i INT PRIMARY KEY,
    j INT
) ENGINE = InnoDB;

CREATE TABLE fc2 (
    m INT PRIMARY KEY,
    n INT,
    FOREIGN KEY ni (n) REFERENCES fc1 (i)
        ON DELETE CASCADE
) ENGINE = InnoDB;
```

Suppose that the slave does not have `InnoDB` support enabled. If this is the case, then the tables on the slave are created, but they use the `MyISAM` storage engine, and the `FOREIGN KEY` option is ignored. Now we insert some rows into the tables on the master:

```
master> INSERT INTO fc1 VALUES (1, 1), (2, 2);
Query OK, 2 rows affected (0.09 sec)
Records: 2  Duplicates: 0  Warnings: 0

master> INSERT INTO fc2 VALUES (1, 1), (2, 2), (3, 1);
Query OK, 3 rows affected (0.19 sec)
Records: 3  Duplicates: 0  Warnings: 0
```

At this point, on both the master and the slave, table `fc1` contains 2 rows, and table `fc2` contains 3 rows, as shown here:

```
master> SELECT * FROM fc1;
+---+------+
| i | j    |
+---+------+
| 1 |    1 |
| 2 |    2 |
+---+------+
2 rows in set (0.00 sec)

master> SELECT * FROM fc2;
+---+------+
| m | n    |
+---+------+
| 1 |    1 |
| 2 |    2 |
| 3 |    1 |
+---+------+
3 rows in set (0.00 sec)

slave> SELECT * FROM fc1;
+---+------+
| i | j    |
+---+------+
| 1 |    1 |
| 2 |    2 |
+---+------+
2 rows in set (0.00 sec)

slave> SELECT * FROM fc2;
+---+------+
| m | n    |
```

```
+---+------+
| 1 |    1 |
| 2 |    2 |
| 3 |    1 |
+---+------+
3 rows in set (0.00 sec)
```

Now suppose that you perform the following `DELETE` statement on the master:

```
master> DELETE FROM fc1 WHERE i=1;
Query OK, 1 row affected (0.09 sec)
```

Due to the cascade, table `fc2` on the master now contains only 1 row:

```
master> SELECT * FROM fc2;
+---+---+
| m | n |
+---+---+
| 2 | 2 |
+---+---+
1 row in set (0.00 sec)
```

However, the cascade does not propagate on the slave because on the slave the `DELETE` for `fc1` deletes no rows from `fc2`. The slave's copy of `fc2` still contains all of the rows that were originally inserted:

```
slave> SELECT * FROM fc2;
+---+---+
| m | n |
+---+---+
| 1 | 1 |
| 3 | 1 |
| 2 | 2 |
+---+---+
3 rows in set (0.00 sec)
```

This difference is due to the fact that the cascading deletes are handled internally by the `InnoDB` storage engine, which means that none of the changes are logged.

15.19 InnoDB memcached Plugin

The `InnoDB memcached` plugin (`daemon_memcached`) provides an integrated `memcached` daemon that automatically stores and retrieves data from `InnoDB` tables, turning the MySQL server into a fast "key-value store". Instead of formulating queries in SQL, you can use simple `get`, `set`, and `incr` operations that avoid the performance overhead associated with SQL parsing and constructing a query optimization plan. You can also access the same `InnoDB` tables through SQL for convenience, complex queries, bulk operations, and other strengths of traditional database software.

This "NoSQL-style" interface uses the `memcached` API to speed up database operations, letting `InnoDB` handle memory caching using its buffer pool mechanism. Data modified through `memcached` operations such as `add`, `set`, and `incr` are stored to disk, in `InnoDB` tables. The combination of `memcached` simplicity and `InnoDB` reliability and consistency provides users with the best of both worlds, as explained in Section 15.19.1, "Benefits of the InnoDB memcached Plugin". For an architectural overview, see Section 15.19.2, "InnoDB memcached Architecture".

15.19.1 Benefits of the InnoDB memcached Plugin

This section outlines advantages the `daemon_memcached` plugin. The combination of `InnoDB` tables and `memcached` offers advantages over using either by themselves.

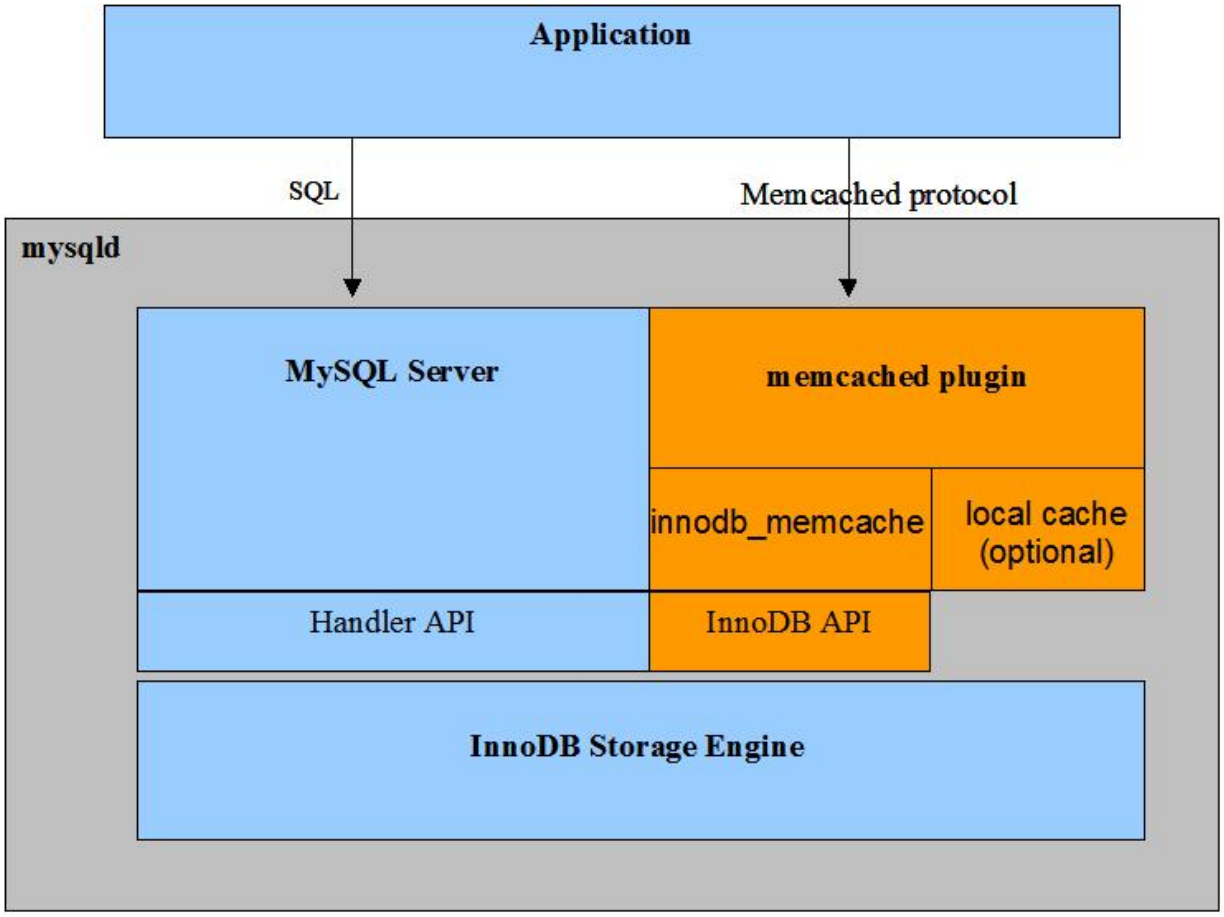
- Direct access to the `InnoDB` storage engine avoids the parsing and planning overhead of SQL.

- Running `memcached` in the same process space as the MySQL server avoids the network overhead of passing requests back and forth.
- Data written using the `memcached` protocol is transparently written to an `InnoDB` table, without going through the MySQL SQL layer. You can control frequency of writes to achieve higher raw performance when updating non-critical data.
- Data requested through the `memcached` protocol is transparently queried from an `InnoDB` table, without going through the MySQL SQL layer.
- Subsequent requests for the same data is served from the `InnoDB` buffer pool. The buffer pool handles the in-memory caching. You can tune performance of data-intensive operations using `InnoDB` configuration options.
- Data can be unstructured or structured, depending on the type of application. You can create a new table for data, or use existing tables.
- `InnoDB` can handle composing and decomposing multiple column values into a single `memcached` item value, reducing the amount of string parsing and concatenation required in your application. For example, you can store the string value `2|4|6|8` in the `memcached` cache, and have `InnoDB` split the value based on a separator character, then store the result in four numeric columns.
- The transfer between memory and disk is handled automatically, simplifying application logic.
- Data is stored in a MySQL database to protect against crashes, outages, and corruption.
- You can access the underlying `InnoDB` table through SQL for reporting, analysis, ad hoc queries, bulk loading, multi-step transactional computations, set operations such as union and intersection, and other operations suited to the expressiveness and flexibility of SQL.
- You can ensure high availability by using the `daemon_memcached` plugin on a master server in combination with MySQL replication.
- The integration of `memcached` with MySQL provides a way to make in-memory data persistent, so you can use it for more significant kinds of data. You can use more `add`, `incr`, and similar write operations in your application without concern that data could be lost. You can stop and start the `memcached` server without losing updates made to cached data. To guard against unexpected outages, you can take advantage of `InnoDB` crash recovery, replication, and backup capabilities.
- The way `InnoDB` does fast primary key lookups is a natural fit for `memcached` single-item queries. The direct, low-level database access path used by the `daemon_memcached` plugin is much more efficient for key-value lookups than equivalent SQL queries.
- The serialization features of `memcached`, which can turn complex data structures, binary files, or even code blocks into storeable strings, offer a simple way to get such objects into a database.
- Because you can access the underlying data through SQL, you can produce reports, search or update across multiple keys, and call functions such as `AVG()` and `MAX()` on `memcached` data. All of these operations are expensive or complicated using `memcached` by itself.
- You do not need to manually load data into `memcached` at startup. As particular keys are requested by an application, values are retrieved from the database automatically, and cached in memory using the `InnoDB` buffer pool.
- Because `memcached` consumes relatively little CPU, and its memory footprint is easy to control, it can run comfortably alongside a MySQL instance on the same system.
- Because data consistency is enforced by mechanisms used for regular `InnoDB` tables, you do not have to worry about stale `memcached` data or fallback logic to query the database in the case of a missing key.

15.19.2 InnoDB memcached Architecture

The `InnoDB` `memcached` plugin implements `memcached` as a MySQL plugin daemon that accesses the `InnoDB` storage engine directly, bypassing the MySQL SQL layer.

The following diagram illustrates how an application accesses data through the `daemon_memcached` plugin, compared with SQL.



Features of the `daemon_memcached` plugin:

- `memcached` as a daemon plugin of `mysqld`. Both `mysqld` and `memcached` run in the same process space, with very low latency access to data.
- Direct access to `InnoDB` tables, bypassing the SQL parser, the optimizer, and even the Handler API layer.
- Standard `memcached` protocols, including the text-based protocol and the binary protocol. The `daemon_memcached` plugin passes all 55 compatibility tests of the `memcapable` command.
- Multi-column support. You can map multiple columns into the “value” part of the key/value store, with column values delimited by a user-specified separator character.
- By default, the `memcached` protocol is used to read and write data directly to `InnoDB`, letting MySQL manage in-memory caching using the `InnoDB` buffer pool. The default settings represent a combination of high reliability and the fewest surprises for database applications. For example, default settings avoid uncommitted data on the database side, or stale data returned for `memcached` `get` requests.
- Advanced users can configure the system as a traditional `memcached` server, with all data cached only in the `memcached` engine (memory caching), or use a combination of the “`memcached` engine” (memory caching) and the `InnoDB` `memcached` engine (`InnoDB` as back-end persistent storage).
- Control over how often data is passed back and forth between `InnoDB` and `memcached` operations through the `innodb_api_bk_commit_interval`, `daemon_memcached_r_batch_size`, and

`daemon_memcached_w_batch_size` configuration options. Batch size options default to a value of 1 for maximum reliability.

- The ability to specify `memcached` options through the `daemon_memcached_option` configuration parameter. For example, you can change the port that `memcached` listens on, reduce the maximum number of simultaneous connections, change the maximum memory size for a key/value pair, or enable debugging messages for the error log.

- The `innodb_api_trx_level` configuration option controls the transaction isolation level on queries processed by `memcached`. Although `memcached` has no concept of transactions, you can use this option to control how soon `memcached` sees changes caused by SQL statements issued on the table used by the `daemon_memcached` plugin. By default, `innodb_api_trx_level` is set to `READ UNCOMMITTED`.

- The `innodb_api_enable_mdl` option can be used to lock the table at the MySQL level, so that the mapped table cannot be dropped or altered by DDL through the SQL interface. Without the lock, the table can be dropped from the MySQL layer, but kept in `InnoDB` storage until `memcached` or some other user stops using it. “MDL” stands for “metadata locking”.

Differences Between InnoDB memcached and Traditional memcached

You may already be familiar with using `memcached` with MySQL, as described in Section 17.2, “Using MySQL with `memcached`”. This section describes how features of the integrated `InnoDB` `memcached` plugin differ from traditional `memcached`.

- Installation: The `memcached` library comes with the MySQL server, making installation and setup relatively easy. Installation involves running the `innodb_memcached_config.sql` script to create a `demo_test` table for `memcached` to use, issuing an `INSTALL PLUGIN` statement to enable the `daemon_memcached` plugin, and adding desired `memcached` options to a MySQL configuration file or startup script. You might still install the traditional `memcached` distribution for additional utilities such as `memcp`, `memcat`, and `memcapable`.

  For comparison with traditional `memcached`, see Section 17.2.1, “Installing `memcached`”.

- Deployment: With traditional `memcached`, it is typical to run large numbers of low-capacity `memcached` servers. A typical deployment of the `daemon_memcached` plugin, however, involves a smaller number of moderate or high-powered servers that are already running MySQL. The benefit of this configuration is in improving efficiency of individual database servers rather than exploiting unused memory or distributing lookups across large numbers of servers. In the default configuration, very little memory is used for `memcached`, and in-memory lookups are served from the `InnoDB` buffer pool, which automatically caches the most recently and frequently used data. As with a traditional MySQL server instance, keep the value of the `innodb_buffer_pool_size` configuration option as high as practical (without causing paging at the OS level), so that as much work as possible is performed in memory.

  For comparison with traditional `memcached`, see Section 17.2.2.2, “`memcached` Deployment”.

- Expiry: By default (that is, using the `innodb_only` caching policy), the latest data from the `InnoDB` table is always returned, so the expiry options have no practical effect. If you change the caching policy to `caching` or `cache-only`, the expiry options work as usual, but requested data might be stale if it is updated in the underlying table before it expires from the memory cache.

  For comparison with traditional `memcached`, see Section 17.2.2.4, “Data Expiry”.

- Namespaces: `memcached` is like a large directory where you give files elaborate names with prefixes and suffixes to keep the files from conflicting. The `daemon_memcached` plugin lets you use similar naming conventions for keys, with one addition. Key names in the format `@@table_id.key.table_id` are decoded to reference a specific a table, using mapping data from the `innodb_memcache.containers` table. The `key` is looked up in or written to the specified table.

The `@@` notation only works for individual calls to `get`, `add`, and `set` functions, but not others such as `incr` or `delete`. To designate a default table for subsequent `memcached` operations within a session, perform a `get` request using the `@@` notation with a *`table_id`*, but without the key portion. For example:

```
get @@table_id
```

Subsequent `get`, `set`, `incr`, `delete`, and other operations use the table designated by *`table_id`* in the `innodb_memcache.containers.name` column.

For comparison with traditional `memcached`, see Section 17.2.2.3, “Using Namespaces”.

- Hashing and distribution: The default configuration, which uses the `innodb_only` caching policy, is suitable for a traditional deployment configuration where all data is available on all servers, such as a set of replication slave servers.

  If you physically divide data, as in a sharded configuration, you can split data across several machines running the `daemon_memcached` plugin, and use the traditional `memcached` hashing mechanism to route requests to a particular machine. On the MySQL side, you would typically let all data be inserted by `add` requests to `memcached` so that appropriate values are stored in the database on the appropriate server.

  For comparison with traditional `memcached`, see Section 17.2.2.5, “`memcached` Hashing/Distribution Types”.

- Memory usage: By default (with the `innodb_only` caching policy), the `memcached` protocol passes information back and forth with `InnoDB` tables, and the `InnoDB` buffer pool handles in-memory lookups instead of `memcached` memory usage growing and shrinking. Relatively little memory is used on the `memcached` side.

  If you switch the caching policy to `caching` or `cache-only`, the normal rules of `memcached` memory usage apply. Memory for `memcached` data values is allocated in terms of “slabs”. You can control slab size and maximum memory used for `memcached`.

  Either way, you can monitor and troubleshoot the `daemon_memcached` plugin using the familiar statistics system, accessed through the standard protocol, over a `telnet` session, for example. Extra utilities are not included with the `daemon_memcached` plugin. You can use the `memcached-tool` script to install a full `memcached` distribution.

  For comparison with traditional `memcached`, see Section 17.2.2.7, “Memory Allocation within `memcached`”.

- Thread usage: MySQL threads and `memcached` threads co-exist on the same server. Limits imposed on threads by the operating system apply to the total number of threads.

  For comparison with traditional `memcached`, see Section 17.2.2.8, “`memcached` Thread Support”.

- Log usage: Because the `memcached` daemon is run alongside the MySQL server and writes to `stderr`, the `-v`, `-vv`, and `-vvv` options for logging write output to the MySQL error log.

  For comparison with traditional `memcached`, see Section 17.2.2.9, “`memcached` Logs”.

- `memcached` operations: Familiar `memcached` operations such as `get`, `set`, `add`, and `delete` are available. Serialization (that is, the exact string format representing complex data structures) depends on the language interface.

  For comparison with traditional `memcached`, see Section 17.2.3.1, “Basic `memcached` Operations”.

- Using `memcached` as a MySQL front end: This is the primary purpose of the `InnoDB` `memcached` plugin. An integrated `memcached` daemon improves application performance, and having `InnoDB` handle data transfers between memory and disk simplifies application logic.

For comparison with traditional `memcached`, see Section 17.2.3.2, “Using `memcached` as a MySQL Caching Layer”.

- Utilities: The MySQL server includes the `libmemcached` library but not additional command-line utilities. To use commands such as `memcp`, `memcat`, and `memcapable` commands, install a full `memcached` distribution. When `memrm` and `memflush` remove items from the cache, the items are also removed from the underlying `InnoDB` table.

  For comparison with traditional `memcached`, see `libmemcached` Command-Line Utilities.

- Programming interfaces: You can access the MySQL server through the `daemon_memcached` plugin using all supported languages: C and C++, Java, Perl, Python, PHP, and Ruby. Specify the server hostname and port as with a traditional `memcached` server. By default, the `daemon_memcached` plugin listens on port `11211`. You can use both the text and binary protocols. You can customize the behavior of `memcached` functions at runtime. Serialization (that is, the exact string format representing complex data structures) depends on the language interface.

  For comparison with traditional `memcached`, see Section 17.2.3, “Developing a `memcached` Application”.

- Frequently asked questions: MySQL has an extensive FAQ for traditional `memcached`. The FAQ is mostly applicable, except that using `InnoDB` tables as a storage medium for `memcached` data means that you can use `memcached` for more write-intensive applications than before, rather than as a read-only cache.

  See Section 17.2.5, “`memcached` FAQ”.

15.19.3 Setting Up the InnoDB memcached Plugin

This section describes how to set up the `daemon_memcached` plugin on a MySQL server. Because the `memcached` daemon is tightly integrated with the MySQL server to avoid network traffic and minimize latency, you perform this process on each MySQL instance that uses this feature.

Note

Before setting up the `daemon_memcached` plugin, consult Section 15.19.5, “Security Considerations for the InnoDB memcached Plugin” to understand the security procedures required to prevent unauthorized access.

Prerequisites

- The `daemon_memcached` plugin is only supported on Linux, Solaris, and OS X platforms. Other operating systems are not supported.
- When building MySQL from source, you must build with `-DWITH_INNODB_MEMCACHED=ON`. This build option generates two shared libraries in the MySQL plugin directory (`plugin_dir`) that are required to run the `daemon_memcached` plugin:
  - `libmemcached.so`: the `memcached` daemon plugin to MySQL.
  - `innodb_engine.so`: an `InnoDB` API plugin to `memcached`.
- `libevent` must be installed.
  - If you did not build MySQL from source, the `libevent` library is not included in your installation. Use the installation method for your operating system to install `libevent` 1.4.12 or later. For example, depending on the operating system, you might use `apt-get`, `yum`, or `port install`. For example, on Ubuntu Linux, use:

    ```
    sudo apt-get install libevent-dev
    ```

- If you installed MySQL from a source code release, `libevent` 1.4.12 is bundled with the package and is located at the top level of the MySQL source code directory. If you use the bundled version of `libevent`, no action is required. If you want to use a local system version of `libevent`, you must build MySQL with the `-DWITH_LIBEVENT` build option set to `system` or `yes`.

Installing and Configuring the InnoDB memcached Plugin

1. Configure the `daemon_memcached` plugin so it can interact with `InnoDB` tables by running the `innodb_memcached_config.sql` configuration script, which is located in *`MYSQL_HOME`*`/share`. This script installs the `innodb_memcache` database with three required tables (`cache_policies`, `config_options`, and `containers`). It also installs the `demo_test` sample table in the `test` database.

```
mysql> source MYSQL_HOME/share/innodb_memcached_config.sql
```

Running the `innodb_memcached_config.sql` script is a one-time operation. The tables remain in place if you later uninstall and re-install the `daemon_memcached` plugin.

```
mysql> USE innodb_memcache;
mysql> SHOW TABLES;
+---------------------------+
| Tables_in_innodb_memcache |
+---------------------------+
| cache_policies            |
| config_options            |
| containers                |
+---------------------------+

mysql> USE test;
mysql> SHOW TABLES;
+----------------+
| Tables_in_test |
+----------------+
| demo_test      |
+----------------+
```

Of these tables, the `innodb_memcache.containers` table is the most important. Entries in the `containers` table provide a mapping to `InnoDB` table columns. Each `InnoDB` table used with the `daemon_memcached` plugin requires an entry in the `containers` table.

The `innodb_memcached_config.sql` script inserts a single entry in the `containers` table that provides a mapping for the `demo_test` table. It also inserts a single row of data into the `demo_test` table. This data allows you to immediately verify the installation after the setup is completed.

```
mysql> SELECT * FROM innodb_memcache.containers\G
*************************** 1. row ***************************
                name: aaa
           db_schema: test
            db_table: demo_test
         key_columns: c1
       value_columns: c2
               flags: c3
          cas_column: c4
   expire_time_column: c5
unique_idx_name_on_key: PRIMARY

mysql> SELECT * FROM test.demo_test;
+----+-----------------+------+------+------+
| c1 | c2              | c3   | c4   | c5   |
+----+-----------------+------+------+------+
| AA | HELLO, HELLO    |    8 |    0 |    0 |
+----+-----------------+------+------+------+
```

For more information about `innodb_memcache` tables and the `demo_test` sample table, see Section 15.19.8, “InnoDB memcached Plugin Internals”.

2. Activate the `daemon_memcached` plugin by running the `INSTALL PLUGIN` statement:

```
mysql> INSTALL PLUGIN daemon_memcached soname "libmemcached.so";
```

Once the plugin is installed, it is automatically activated each time the MySQL server is restarted.

Verifying the InnoDB and memcached Setup

To verify the `daemon_memcached` plugin setup, use a `telnet` session to issue `memcached` commands. By default, the `memcached` daemon listens on port 11211.

1. Retrieve data from the `test.demo_test` table. The single row of data in the `demo_test` table has a key value of `AA`.

```
telnet localhost 11211
Trying 127.0.0.1...
Connected to localhost.
Escape character is '^]'.
get AA
VALUE AA 8 12
HELLO, HELLO
END
```

2. Insert data using a `set` command.

```
set BB 10 0 16
GOODBYE, GOODBYE
STORED
```

where:

- `set` is the command to store a value
- `BB` is the key
- `10` is a flag for the operation; ignored by `memcached` but may be used by the client to indicate any type of information; specify `0` if unused
- `0` is the expiration time (TTL); specify `0` if unused
- `16` is the length of the supplied value block in bytes
- `GOODBYE, GOODBYE` is the value that is stored

3. Verify that the data inserted is stored in MySQL by connecting to the MySQL server and querying the `test.demo_test` table.

```
mysql> SELECT * FROM test.demo_test;
+----+------------------+------+------+------+
| c1 | c2               | c3   | c4   | c5   |
+----+------------------+------+------+------+
| AA | HELLO, HELLO     |    8 |    0 |    0 |
| BB | GOODBYE, GOODBYE |   10 |    1 |    0 |
+----+------------------+------+------+------+
```

4. Return to the telnet session and retrieve the data that you inserted earlier using key `BB`.

```
get BB
VALUE BB 10 16
GOODBYE, GOODBYE
```

```
END
quit
```

If you shut down the MySQL server, which also shuts off the integrated `memcached` server, further attempts to access the `memcached` data fail with a connection error. Normally, the `memcached` data also disappears at this point, and you would require application logic to load the data back into memory when `memcached` is restarted. However, the `InnoDB` `memcached` plugin automates this process for you.

When you restart MySQL, `get` operations once again return the key/value pairs you stored in the earlier `memcached` session. When a key is requested and the associated value is not already in the memory cache, the value is automatically queried from the MySQL `test.demo_test` table.

Creating a New Table and Column Mapping

This example shows how to setup your own `InnoDB` table with the `daemon_memcached` plugin.

1. Create an `InnoDB` table. The table must have a key column with a unique index. The key column of the city table is `city_id`, which is defined as the primary key. The table must also include columns for `flags`, `cas`, and `expiry` values. There may be one or more value columns. The `city` table has three value columns (`name`, `state`, `country`).

   Note

   There is no special requirement with respect to column names as along as a valid mapping is added to the `innodb_memcache.containers` table.

   ```
   mysql> CREATE TABLE city (
          city_id VARCHAR(32),
          name VARCHAR(1024),
          state VARCHAR(1024),
          country VARCHAR(1024),
          flags INT,
          cas BIGINT UNSIGNED,
          expiry INT,
          primary key(city_id)
          ) ENGINE=InnoDB;
   ```

2. Add an entry to the `innodb_memcache.containers` table so that the `daemon_memcached` plugin knows how to access the `InnoDB` table. The entry must satisfy the `innodb_memcache.containers` table definition. For a description of each field, see Section 15.19.8, “InnoDB memcached Plugin Internals”.

   ```
   mysql> DESCRIBE innodb_memcache.containers;
   +------------------------+--------------+------+-----+---------+-------+
   | Field                  | Type         | Null | Key | Default | Extra |
   +------------------------+--------------+------+-----+---------+-------+
   | name                   | varchar(50)  | NO   | PRI | NULL    |       |
   | db_schema              | varchar(250) | NO   |     | NULL    |       |
   | db_table               | varchar(250) | NO   |     | NULL    |       |
   | key_columns            | varchar(250) | NO   |     | NULL    |       |
   | value_columns          | varchar(250) | YES  |     | NULL    |       |
   | flags                  | varchar(250) | NO   |     | 0       |       |
   | cas_column             | varchar(250) | YES  |     | NULL    |       |
   | expire_time_column     | varchar(250) | YES  |     | NULL    |       |
   | unique_idx_name_on_key | varchar(250) | NO   |     | NULL    |       |
   +------------------------+--------------+------+-----+---------+-------+
   ```

   The `innodb_memcache.containers` table entry for the city table is defined as:

   ```
   mysql> INSERT INTO `innodb_memcache`.`containers` (
          `name`, `db_schema`, `db_table`, `key_columns`, `value_columns`,
          `flags`, `cas_column`, `expire_time_column`, `unique_idx_name_on_key`)
          VALUES ('default', 'test', 'city', 'city_id', 'name|state|country',
          'flags','cas','expiry','PRIMARY');
   ```

- `default` is specified for the `containers.name` column to configure the `city` table as the default `InnoDB` table to be used with the `daemon_memcached` plugin.
- Multiple `InnoDB` table columns (`name`, `state`, `country`) are mapped to `containers.value_columns` using a "|" delimiter.
- The `flags`, `cas_column`, and `expire_time_column` fields of the `innodb_memcache.containers` table are typically not significant in applications using the `daemon_memcached` plugin. However, a designated `InnoDB` table column is required for each. When inserting data, specify `0` for these columns if they are unused.

3. After updating the `innodb_memcache.containers` table, restart the `daemon_memcache` plugin to apply the changes.

```
mysql> UNINSTALL PLUGIN daemon_memcached;

mysql> INSTALL PLUGIN daemon_memcached soname "libmemcached.so";
```

4. Using telnet, insert data into the `city` table using a `memcached set` command.

```
telnet localhost 11211
Trying 127.0.0.1...
Connected to localhost.
Escape character is '^]'.
set B 0 0 22
BANGALORE|BANGALORE|IN
STORED
```

5. Using MySQL, query the `test.city` table to verify that the data you inserted was stored.

```
mysql> SELECT * FROM test.city;
+---------+-----------+-----------+---------+-------+------+--------+
| city_id | name      | state     | country | flags | cas  | expiry |
+---------+-----------+-----------+---------+-------+------+--------+
| B       | BANGALORE | BANGALORE | IN      |     0 |    3 |      0 |
+---------+-----------+-----------+---------+-------+------+--------+
```

6. Using MySQL, insert additional data into the `test.city` table.

```
mysql> INSERT INTO city VALUES ('C','CHENNAI','TAMIL NADU','IN', 0, 0 ,0);
mysql> INSERT INTO city VALUES ('D','DELHI','DELHI','IN', 0, 0, 0);
mysql> INSERT INTO city VALUES ('H','HYDERABAD','TELANGANA','IN', 0, 0, 0);
mysql> INSERT INTO city VALUES ('M','MUMBAI','MAHARASHTRA','IN', 0, 0, 0);
```

Note

It is recommended that you specify a value of `0` for the `flags`, `cas_column`, and `expire_time_column` fields if they are unused.

7. Using telnet, issue a `memcached get` command to retrieve data you inserted using MySQL.

```
get H
VALUE H 0 22
HYDERABAD|TELANGANA|IN
END
```

Configuring the InnoDB memcached Plugin

Traditional `memcached` configuration options may be specified in a MySQL configuration file or a `mysqld` startup string, encoded in the argument of the `daemon_memcached_option` configuration parameter. `memcached` configuration options take effect when the plugin is loaded, which occurs each time the MySQL server is started.

For example, to make `memcached` listen on port 11222 instead of the default port 11211, specify `-p11222` as an argument of the `daemon_memcached_option` configuration option:

```
mysqld .... --daemon_memcached_option="-p11222"
```

Other `memcached` options can be encoded in the `daemon_memcached_option` string. For example, you can specify options to reduce the maximum number of simultaneous connections, change the maximum memory size for a key/value pair, or enable debugging messages for the error log, and so on.

There are also configuration options specific to the `daemon_memcached` plugin. These include:

- `daemon_memcached_engine_lib_name`: Specifies the shared library that implements the `InnoDB memcached` plugin. The default setting is `innodb_engine.so`.
- `daemon_memcached_engine_lib_path`: The path of the directory containing the shared library that implements the `InnoDB memcached` plugin. The default is NULL, representing the plugin directory.
- `daemon_memcached_r_batch_size`: Defines the batch commit size for read operations (`get`). It specifies the number of `memcached` read operations after which a commit occurs. `daemon_memcached_r_batch_size` is set to 1 by default so that every `get` request accesses the most recently committed data in the `InnoDB` table, whether the data was updated through `memcached` or by SQL. When the value is greater than 1, the counter for read operations is incremented with each `get` call. A `flush_all` call resets both read and write counters.
- `daemon_memcached_w_batch_size`: Defines the batch commit size for write operations (`set`, `replace`, `append`, `prepend`, `incr`, `decr`, and so on). `daemon_memcached_w_batch_size` is set to 1 by default so that no uncommitted data is lost in case of an outage, and so that SQL queries on the underlying table access the most recent data. When the value is greater than 1, the counter for write operations is incremented for each `add`, `set`, `incr`, `decr`, and `delete` call. A `flush_all` call resets both read and write counters.

By default, you do not need to modify `daemon_memcached_engine_lib_name` or `daemon_memcached_engine_lib_path`. You might configure these options if, for example, you want to use a different storage engine for `memcached` (such as the NDB `memcached` engine).

`daemon_memcached` plugin configuration parameters may be specified in the MySQL configuration file or in a `mysqld` startup string. They take effect when you load the `daemon_memcached` plugin.

When making changes to `daemon_memcached` plugin configuration, reload the plugin to apply the changes. To do so, issue the following statements:

```
mysql> UNINSTALL PLUGIN daemon_memcached;

mysql> INSTALL PLUGIN daemon_memcached soname "libmemcached.so";
```

Configuration settings, required tables, and data are preserved when the plugin is restarted.

For additional information about enabling and disabling plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

15.19.4 InnoDB memcached Multiple get and Range Query Support

The `daemon_memcached` plugin supports multiple get operations (fetching multiple key/value pairs in a single `memcached` query) and range queries.

Multiple get Operations

The ability to fetch multiple key/value pairs in a single `memcached` query improves read performance by reducing communication traffic between the client and server. For `InnoDB`, it means fewer transactions and open-table operations.

The following example demonstrates multiple-get support. The example uses the `test.city` table described in Creating a New Table and Column Mapping.

```
mysql> USE test;
mysql> SELECT * FROM test.city;
+---------+-----------+-------------+---------+-------+------+--------+
| city_id | name      | state       | country | flags | cas  | expiry |
+---------+-----------+-------------+---------+-------+------+--------+
| B       | BANGALORE | BANGALORE   | IN      |     0 |    1 |      0 |
| C       | CHENNAI   | TAMIL NADU  | IN      |     0 |    0 |      0 |
| D       | DELHI     | DELHI       | IN      |     0 |    0 |      0 |
| H       | HYDERABAD | TELANGANA   | IN      |     0 |    0 |      0 |
| M       | MUMBAI    | MAHARASHTRA | IN      |     0 |    0 |      0 |
+---------+-----------+-------------+---------+-------+------+--------+
```

Run a `get` command to retrieve all values from the `city` table. The results are returned in a key/value pair sequence.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
get B C D H M
VALUE B 0 22
BANGALORE|BANGALORE|IN
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
VALUE M 0 21
MUMBAI|MAHARASHTRA|IN
END
```

When retrieving multiple values in a single `get` command, you can switch tables (using *`@@containers.name`* notation) to retrieve the value for the first key, but you cannot switch tables for subsequent keys. For example, the table switch in this example is valid:

```
get @@aaa.AA BB
VALUE @@aaa.AA 8 12
HELLO, HELLO
VALUE BB 10 16
GOODBYE, GOODBYE
END
```

Attempting to switch tables again in the same `get` command to retrieve a key value from a different table is not supported.

Range Queries

For range queries, the `daemon_memcached` plugin supports the following comparison operators: `<`, `>`, `<=`, `>=`. An operator must be preceded by an `@` symbol. When a range query finds multiple matching key/value pairs, results are returned in a key/value pair sequence.

The following examples demonstrate range query support. The examples use the `test.city` table described in Creating a New Table and Column Mapping.

```
mysql> SELECT * FROM test.city;
+---------+-----------+-------------+---------+-------+------+--------+
| city_id | name      | state       | country | flags | cas  | expiry |
+---------+-----------+-------------+---------+-------+------+--------+
| B       | BANGALORE | BANGALORE   | IN      |     0 |    1 |      0 |
| C       | CHENNAI   | TAMIL NADU  | IN      |     0 |    0 |      0 |
| D       | DELHI     | DELHI       | IN      |     0 |    0 |      0 |
```

```
| H       | HYDERABAD | TELANGANA   | IN      |     0 |    0 |      0 |
| M       | MUMBAI    | MAHARASHTRA | IN      |     0 |    0 |      0 |
+---------+-----------+-------------+---------+-------+------+--------+
```

Open a telnet session:

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
```

To get all values greater than `B`, enter `get @>B`:

```
get @>B
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
VALUE M 0 21
MUMBAI|MAHARASHTRA|IN
END
```

To get all values less than `M`, enter `get @<M`:

```
get @<M
VALUE B 0 22
BANGALORE|BANGALORE|IN
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
END
```

To get all values less than and including `M`, enter `get @<=M`:

```
get @<=M
VALUE B 0 22
BANGALORE|BANGALORE|IN
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
VALUE M 0 21
MUMBAI|MAHARASHTRA|IN
```

To get values greater than `B` but less than `M`, enter `get @>B@<M`:

```
get @>B@<M
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
END
```

A maximum of two comparison operators can be parsed, one being either a 'less than' (`@<`) or 'less than or equal to' (`@<=`) operator, and the other being either a 'greater than' (`@>`) or 'greater than or

equal to' (`@>=`) operator. Any additional operators are assumed to be part of the key. For example, if you issue a `get` command with three operators, the third operator (`@>C`) is treated as part of the key, and the `get` command searches for values smaller than `M` and greater than `B@>C`.

```
get @<M@>B@>C
VALUE C 0 21
CHENNAI|TAMIL NADU|IN
VALUE D 0 14
DELHI|DELHI|IN
VALUE H 0 22
HYDERABAD|TELANGANA|IN
```

15.19.5 Security Considerations for the InnoDB memcached Plugin

Caution

Consult this section before deploying the `daemon_memcached` plugin on a production server, or even on a test server if the MySQL instance contains sensitive data.

Because `memcached` does not use an authentication mechanism by default, and the optional SASL authentication is not as strong as traditional DBMS security measures, only keep non-sensitive data in the MySQL instance that uses the `daemon_memcached` plugin, and wall off any servers that use this configuration from potential intruders. Do not allow `memcached` access to these servers from the Internet; only allow access from within a firewalled intranet, ideally from a subnet whose membership you can restrict.

Password-Protecting memcached Using SASL

SASL support provides the capability to protect your MySQL database from unauthenticated access through `memcached` clients. This section explains how to enable SASL with the `daemon_memcached` plugin. The steps are almost identical to those performed to enabled SASL for a traditional `memcached` server.

SASL stands for “Simple Authentication and Security Layer”, a standard for adding authentication support to connection-based protocols. `memcached` added SASL support in version 1.4.3.

SASL authentication is only supported with the binary protocol.

`memcached` clients are only able to access `InnoDB` tables that are registered in the `innodb_memcache.containers` table. Even though a DBA can place access restrictions on such tables, access through `memcached` applications cannot be controlled. For this reason, SASL support is provided to control access to `InnoDB` tables associated with the `daemon_memcached` plugin.

The following section shows how to build, enable, and test an SASL-enabled `daemon_memcached` plugin.

Building and Enabling SASL with the InnoDB memcached Plugin

By default, an SASL-enabled `daemon_memcached` plugin is not included in MySQL release packages, since an SASL-enabled `daemon_memcached` plugin requires building `memcached` with SASL libraries. To enable SASL support, download the MySQL source and rebuild the `daemon_memcached` plugin after downloading the SASL libraries:

1. Install the SASL development and utility libraries. For example, on Ubuntu, use `apt-get` to obtain the libraries:

```
sudo apt-get -f install libsasl2-2 sasl2-bin libsasl2-2 libsasl2-dev libsasl2-modules
```

2. Build the `daemon_memcached` plugin shared libraries with SASL capability by adding `ENABLE_MEMCACHED_SASL=1` to your `cmake` options. `memcached` also provides *simple cleartext password support*, which facilitates testing. To enable simple cleartext password support, specify the `ENABLE_MEMCACHED_SASL_PWDB=1` `cmake` option.

   In summary, add following three `cmake` options:

   ```
   cmake ... -DWITH_INNODB_MEMCACHED=1 -DENABLE_MEMCACHED_SASL=1 -DENABLE_MEMCACHED_SASL_PWDB=1
   ```

3. Install the `daemon_memcached` plugin, as described in Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

4. Configure a user name and password file. (This example uses `memcached` simple cleartext password support.)

   a. In a file, create a user named `testname` and define the password as `testpasswd`:

   ```
   echo "testname:testpasswd:::::::" >/home/jy/memcached-sasl-db
   ```

   b. Configure the `MEMCACHED_SASL_PWDB` environment variable to inform `memcached` of the user name and password file:

   ```
   export MEMCACHED_SASL_PWDB=/home/jy/memcached-sasl-db
   ```

   c. Inform `memcached` that a cleartext password is used:

   ```
   echo "mech_list: plain" > /home/jy/work2/msasl/clients/memcached.conf
   export SASL_CONF_PATH=/home/jy/work2/msasl/clients
   ```

5. Enable SASL by restarting the MySQL server with the `memcached -S` option encoded in the `daemon_memcached_option` configuration parameter:

   ```
   mysqld ... --daemon_memcached_option="-S"
   ```

6. To test the setup, use an SASL-enabled client such as SASL-enabled libmemcached.

   ```
   memcp --servers=localhost:11211 --binary  --username=testname
     --password=testpasswd myfile.txt

   memcat --servers=localhost:11211 --binary --username=testname
     --password=testpasswd myfile.txt
   ```

   If you specify an incorrect user name or password, the operation is rejected with a `memcache error AUTHENTICATION FAILURE` message. In this case, examine the cleartext password set in the `memcached-sasl-db` file to verify that the credentials you supplied are correct.

There are other methods to test SASL authentication with `memcached`, but the method described above is the most straightforward.

15.19.6 Writing Applications for the InnoDB memcached Plugin

Typically, writing an application for the `InnoDB` `memcached` plugin involves some degree of rewriting or adapting existing code that uses MySQL or the `memcached` API.

- With the `daemon_memcached` plugin, instead of many traditional `memcached` servers running on low-powered machines, you have the same number of `memcached` servers as MySQL servers, running on relatively high-powered machines with substantial disk storage and memory. You might reuse some existing code that works with the `memcached` API, but adaptation is likely required due to the different server configuration.

- The data stored through the `daemon_memcached` plugin goes into `VARCHAR`, `TEXT`, or `BLOB` columns, and must be converted to do numeric operations. You can perform the conversion on the application side, or by using the `CAST()` function in queries.
- Coming from a database background, you might be used to general-purpose SQL tables with many columns. The tables accessed by `memcached` code likely have only a few or even a single column holding data values.
- You might adapt parts of your application that perform single-row queries, inserts, updates, or deletes, to improve performance in critical sections of code. Both queries (read) and DML (write) operations can be substantially faster when performed through the `InnoDB memcached` interface. The performance improvement for writes is typically greater than the performance improvement for reads, so you might focus on adapting code that performs logging or records interactive choices on a web site.

The following sections explore these points in more detail.

15.19.6.1 Adapting an Existing MySQL Schema for the InnoDB memcached Plugin

Consider these aspects of `memcached` applications when adapting an existing MySQL schema or application to use the `daemon_memcached` plugin:

- `memcached` keys cannot contain spaces or newlines, because these characters are used as separators in the ASCII protocol. If you are using lookup values that contain spaces, transform or hash them into values without spaces before using them as keys in calls to `add()`, `set()`, `get()`, and so on. Although theoretically these characters are allowed in keys in programs that use the binary protocol, you should restrict the characters used in keys to ensure compatibility with a broad range of clients.
- If there is a short numeric primary key column in an `InnoDB` table, use it as the unique lookup key for `memcached` by converting the integer to a string value. If the `memcached` server is used for multiple applications, or with more than one `InnoDB` table, consider modifying the name to ensure that it is unique. For example, prepend the table name, or the database name and the table name, before the numeric value.

Note

The `daemon_memcached` plugin supports inserts and reads on mapped `InnoDB` tables that have an `INTEGER` defined as the primary key.

- You cannot use a partitioned table for data queried or stored using `memcached`.
- The `memcached` protocol passes numeric values around as strings. To store numeric values in the underlying `InnoDB` table, to implement counters that can be used in SQL functions such as `SUM()` or `AVG()`, for example:
  - Use `VARCHAR` columns with enough characters to hold all the digits of the largest expected number (and additional characters if appropriate for the negative sign, decimal point, or both).
  - In any query that performs arithmetic using column values, use the `CAST()` function to convert the values from string to integer, or to some other numeric type. For example:

```
-- Alphabetic entries are returned as zero.
select cast(c2 as unsigned integer) from demo_test;
-- Since there could be numeric values of 0, can't disqualify them.
-- Test the string values to find the ones that are integers, and average only those.
select avg(cast(c2 as unsigned integer)) from demo_test
  where c2 between '0' and '9999999999';
-- Views let you hide the complexity of queries. The results are already converted;
-- no need to repeat conversion functions and WHERE clauses each time.
create view numbers as select c1 key, cast(c2 as unsigned integer) val
  from demo_test where c2 between '0' and '9999999999';
select sum(val) from numbers;
```

Note

Any alphabetic values in the result set are converted into 0 by the call to `CAST()`. When using functions such as `AVG()`, which depend on the number of rows in the result set, include `WHERE` clauses to filter out non-numeric values.

- If the `InnoDB` column used as a key could have values longer than 250 bytes, hash the value to less than 250 bytes.
- To use an existing table with the `daemon_memcached` plugin, define an entry for it in the `innodb_memcache.containers` table. To make that table the default for all `memcached` requests, specify a value of `default` in the `name` column, then restart the MySQL server to make the change take effect. If you use multiple tables for different classes of `memcached` data, set up multiple entries in the `innodb_memcache.containers` table with `name` values of your choice, then issue a `memcached` request in the form of `get @@name` or `set @@name` within the application to specify the table to be used for subsequent `memcached` requests.

  For an example of using a table other than the predefined `test.demo_test` table, see Example 15.13, "Using Your Own Table with an InnoDB memcached Application". For the required table layout, see Section 15.19.8, "InnoDB memcached Plugin Internals".
- To use multiple `InnoDB` table column values with `memcached` key/value pairs, specify column names separated by comma, semicolon, space, or pipe characters in the `value_columns` field of the `innodb_memcache.containers` entry for the `InnoDB` table. For example, specify `col1,col2,col3` or `col1|col2|col3` in the `value_columns` field.

  Concatenate the column values into a single string using the pipe character as a separator before passing the string to `memcached add` or `set` calls. The string is unpacked automatically into the correct column. Each `get` call returns a single string containing the column values that is also delimited by the pipe character. You can unpack the values using the appropriate application language syntax.

Example 15.13 Using Your Own Table with an InnoDB memcached Application

This example shows how to use your own table with a sample Python application that uses `memcached` for data manipulation.

The example assumes that the `daemon_memcached` plugin is installed as described in Section 15.19.3, "Setting Up the InnoDB memcached Plugin". It also assumes that your system is configured to run a Python script that uses the `python-memcache` module.

1. Create the `multicol` table which stores country information including population, area, and driver side data (`'R'` for right and `'L'` for left).

```
mysql> USE test;

mysql> CREATE TABLE `multicol` (
        `country` varchar(128) NOT NULL DEFAULT '',
        `population` varchar(10) DEFAULT NULL,
        `area_sq_km` varchar(9) DEFAULT NULL,
        `drive_side` varchar(1) DEFAULT NULL,
        `c3` int(11) DEFAULT NULL,
        `c4` bigint(20) unsigned DEFAULT NULL,
        `c5` int(11) DEFAULT NULL,
       PRIMARY KEY (`country`)
       ) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4;
```

2. Insert a record into the `innodb_memcache.containers` table so that the `daemon_memcached` plugin can access the `multicol` table.

```
mysql> INSERT INTO innodb_memcache.containers
       (name,db_schema,db_table,key_columns,value_columns,flags,cas_column,
        expire_time_column,unique_idx_name_on_key)
      VALUES
       ('bbb','test','multicol','country','population,area_sq_km,drive_side',
        'c3','c4','c5','PRIMARY');

mysql> COMMIT;
```

- The `innodb_memcache.containers` record for the `multicol` table specifies a `name` value of `'bbb'`, which is the table identifier.

  Note

  If a single `InnoDB` table is used for all `memcached` applications, the `name` value can be set to `default` to avoid using `@@` notation to switch tables.

- The `db_schema` column is set to `test`, which is the name of the database where the `multicol` table resides.
- The `db_table` column is set to `multicol`, which is the name of the `InnoDB` table.
- `key_columns` is set to the unique `country` column. The `country` column is defined as the primary key in the `multicol` table definition.
- Rather than a single `InnoDB` table column to hold a composite data value, data is divided among three table columns (`population`, `area_sq_km`, and `drive_side`). To accommodate multiple value columns, a comma-separated list of columns is specified in the `value_columns` field. The columns defined in the `value_columns` field are the columns used when storing or retrieving values.
- Values for the `flags`, `expire_time`, and `cas_column` fields are based on values used in the `demo.test` sample table. These fields are typically not significant in applications that use the `daemon_memcached` plugin because MySQL keeps data synchronized, and there is no need to worry about data expiring or becoming stale.
- The `unique_idx_name_on_key` field is set to `PRIMARY`, which refers to the primary index defined on the unique `country` column in the `multicol` table.

3. Copy the sample Python application into a file. In this example, the sample script is copied to a file named `multicol.py`.

   The sample Python application inserts data into the `multicol` table and retrieves data for all keys, demonstrating how to access an `InnoDB` table through the `daemon_memcached` plugin.

```
import sys, os
import memcache

def connect_to_memcached():
  memc = memcache.Client(['127.0.0.1:11211'], debug=0);
  print "Connected to memcached."
  return memc

def banner(message):
  print
  print "=" * len(message)
  print message
  print "=" * len(message)

country_data = [
("Canada","34820000","9984670","R"),
("USA","314242000","9826675","R"),
("Ireland","6399152","84421","L"),
("UK","62262000","243610","L"),
("Mexico","113910608","1972550","R"),
```

```
("Denmark","5543453","43094","R"),
("Norway","5002942","385252","R"),
("UAE","8264070","83600","R"),
("India","1210193422","3287263","L"),
("China","1347350000","9640821","R"),
]

def switch_table(memc,table):
  key = "@@" + table
  print "Switching default table to '" + table + "' by issuing GET for '" + key + "'."
  result = memc.get(key)

def insert_country_data(memc):
  banner("Inserting initial data via memcached interface")
  for item in country_data:
    country = item[0]
    population = item[1]
    area = item[2]
    drive_side = item[3]

    key = country
    value = "|".join([population,area,drive_side])
    print "Key = " + key
    print "Value = " + value

    if memc.add(key,value):
      print "Added new key, value pair."
    else:
      print "Updating value for existing key."
      memc.set(key,value)

def query_country_data(memc):
  banner("Retrieving data for all keys (country names)")
  for item in country_data:
    key = item[0]
    result = memc.get(key)
    print "Here is the result retrieved from the database for key " + key + ":"
    print result
    (m_population, m_area, m_drive_side) = result.split("|")
    print "Unpacked population value: " + m_population
    print "Unpacked area value      : " + m_area
    print "Unpacked drive side value: " + m_drive_side

if __name__ == '__main__':

  memc = connect_to_memcached()
  switch_table(memc,"bbb")
  insert_country_data(memc)
  query_country_data(memc)

  sys.exit(0)
```

Sample Python application notes:

- No database authorization is required to run the application, since data manipulation is performed through the `memcached` interface. The only required information is the port number on the local system where the `memcached` daemon listens.
- To make sure the application uses the `multicol` table, the `switch_table()` function is called, which performs a dummy `get` or `set` request using `@@` notation. The `name` value in the request is `bbb`, which is the `multicol` table identifier defined in the `innodb_memcache.containers.name` field.

  A more descriptive `name` value might be used in a real-world application. This example simply illustrates that a table identifier is specified rather than the table name in `get @@...` requests.
- The utility functions used to insert and query data demonstrate how to turn a Python data structure into pipe-separated values for sending data to MySQL with `add` or `set` requests, and

how to unpack the pipe-separated values returned by `get` requests. This extra processing is only required when mapping a single `memcached` value to multiple MySQL table columns.

4. Run the sample Python application.

```
shell> python multicol.py
```

If successful, the sample application returns this output:

```
Connected to memcached.
Switching default table to 'bbb' by issuing GET for '@@bbb'.

==============================================
Inserting initial data via memcached interface
==============================================
Key = Canada
Value = 34820000|9984670|R
Added new key, value pair.
Key = USA
Value = 314242000|9826675|R
Added new key, value pair.
Key = Ireland
Value = 6399152|84421|L
Added new key, value pair.
Key = UK
Value = 62262000|243610|L
Added new key, value pair.
Key = Mexico
Value = 113910608|1972550|R
Added new key, value pair.
Key = Denmark
Value = 5543453|43094|R
Added new key, value pair.
Key = Norway
Value = 5002942|385252|R
Added new key, value pair.
Key = UAE
Value = 8264070|83600|R
Added new key, value pair.
Key = India
Value = 1210193422|3287263|L
Added new key, value pair.
Key = China
Value = 1347350000|9640821|R
Added new key, value pair.

==========================================
Retrieving data for all keys (country names)
==========================================
Here is the result retrieved from the database for key Canada:
34820000|9984670|R
Unpacked population value: 34820000
Unpacked area value      : 9984670
Unpacked drive side value: R
Here is the result retrieved from the database for key USA:
314242000|9826675|R
Unpacked population value: 314242000
Unpacked area value      : 9826675
Unpacked drive side value: R
Here is the result retrieved from the database for key Ireland:
6399152|84421|L
Unpacked population value: 6399152
Unpacked area value      : 84421
Unpacked drive side value: L
Here is the result retrieved from the database for key UK:
62262000|243610|L
Unpacked population value: 62262000
Unpacked area value      : 243610
Unpacked drive side value: L
Here is the result retrieved from the database for key Mexico:
113910608|1972550|R
```

```
Unpacked population value: 113910608
Unpacked area value      : 1972550
Unpacked drive side value: R
Here is the result retrieved from the database for key Denmark:
5543453|43094|R
Unpacked population value: 5543453
Unpacked area value      : 43094
Unpacked drive side value: R
Here is the result retrieved from the database for key Norway:
5002942|385252|R
Unpacked population value: 5002942
Unpacked area value      : 385252
Unpacked drive side value: R
Here is the result retrieved from the database for key UAE:
8264070|83600|R
Unpacked population value: 8264070
Unpacked area value      : 83600
Unpacked drive side value: R
Here is the result retrieved from the database for key India:
1210193422|3287263|L
Unpacked population value: 1210193422
Unpacked area value      : 3287263
Unpacked drive side value: L
Here is the result retrieved from the database for key China:
1347350000|9640821|R
Unpacked population value: 1347350000
Unpacked area value      : 9640821
Unpacked drive side value: R
dtprice@ubuntu:~$
```

5. Query the `innodb_memcache.containers` table to view the record you inserted earlier for the `multicol` table. The first record is the sample entry for the `demo_test` table that is created during the initial `daemon_memcached` plugin setup. The second record is the entry you inserted for the `multicol` table.

```
mysql> SELECT * FROM innodb_memcache.containers\G
*************************** 1. row ***************************
                  name: aaa
             db_schema: test
              db_table: demo_test
           key_columns: c1
         value_columns: c2
                 flags: c3
            cas_column: c4
    expire_time_column: c5
unique_idx_name_on_key: PRIMARY
*************************** 2. row ***************************
                  name: bbb
             db_schema: test
              db_table: multicol
           key_columns: country
         value_columns: population,area_sq_km,drive_side
                 flags: c3
            cas_column: c4
    expire_time_column: c5
unique_idx_name_on_key: PRIMARY
```

6. Query the `multicol` table to view data inserted by the sample Python application. The data is available for MySQL queries, which demonstrates how the same data can be accessed using SQL or through applications (using the appropriate MySQL Connector or API).

```
mysql> SELECT * FROM test.multicol;
+---------+------------+------------+------------+------+------+------+
| country | population | area_sq_km | drive_side | c3   | c4   | c5   |
+---------+------------+------------+------------+------+------+------+
| Canada  | 34820000   | 9984670    | R          |    0 |   11 |    0 |
| China   | 1347350000 | 9640821    | R          |    0 |   20 |    0 |
| Denmark | 5543453    | 43094      | R          |    0 |   16 |    0 |
| India   | 1210193422 | 3287263    | L          |    0 |   19 |    0 |
| Ireland | 6399152    | 84421      | L          |    0 |   13 |    0 |
| Mexico  | 113910608  | 1972550    | R          |    0 |   15 |    0 |
```

```
| Norway  | 5002942    | 385252     | R          |    0 |   17 |    0 |
| UAE     | 8264070    | 83600      | R          |    0 |   18 |    0 |
| UK      | 62262000   | 243610     | L          |    0 |   14 |    0 |
| USA     | 314242000  | 9826675    | R          |    0 |   12 |    0 |
+---------+------------+------------+------------+------+------+------+
```

Note

Always allow sufficient size to hold necessary digits, decimal points, sign characters, leading zeros, and so on when defining the length for columns that are treated as numbers. Too-long values in a string column such as a `VARCHAR` are truncated by removing some characters, which could produce nonsensical numeric values.

7. Optionally, run report-type queries on the `InnoDB` table that stores the `memcached` data.

   You can produce reports through SQL queries, performing calculations and tests across any columns, not just the `country` key column. (Because the following examples use data from only a few countries, the numbers are for illustration purposes only.) The following queries return the average population of countries where people drive on the right, and the average size of countries whose names start with “U”:

```
mysql> SELECT AVG(population) FROM multicol WHERE drive_side = 'R';
+-------------------+
| avg(population)   |
+-------------------+
| 261304724.7142857 |
+-------------------+

mysql> SELECT SUM(area_sq_km) FROM multicol WHERE country LIKE 'U%';
+-----------------+
| sum(area_sq_km) |
+-----------------+
|        10153885 |
+-----------------+
```

   Because the `population` and `area_sq_km` columns store character data rather than strongly typed numeric data, functions such as `AVG()` and `SUM()` work by converting each value to a number first. This approach *does not work* for operators such as `<` or `>`, for example, when comparing character-based values, `9 > 1000`, which is not expected from a clause such as `ORDER BY population DESC`. For the most accurate type treatment, perform queries against views that cast numeric columns to the appropriate types. This technique lets you issue simple `SELECT *` queries from database applications, while ensuring that casting, filtering, and ordering is correct. The following example shows a view that can be queried to find the top three countries in descending order of population, with the results reflecting the latest data in the `multicol` table, and with population and area figures treated as numbers:

```
mysql> CREATE VIEW populous_countries AS
       SELECT
          country,
          cast(population as unsigned integer) population,
          cast(area_sq_km as unsigned integer) area_sq_km,
          drive_side FROM multicol
        ORDER BY CAST(population as unsigned integer) DESC
        LIMIT 3;

mysql> SELECT * FROM populous_countries;
+---------+------------+------------+------------+
| country | population | area_sq_km | drive_side |
+---------+------------+------------+------------+
| China   | 1347350000 |    9640821 | R          |
| India   | 1210193422 |    3287263 | L          |
| USA     |  314242000 |    9826675 | R          |
+---------+------------+------------+------------+
```

```
mysql> DESC populous_countries;
+------------+---------------------+------+-----+---------+-------+
| Field      | Type                | Null | Key | Default | Extra |
+------------+---------------------+------+-----+---------+-------+
| country    | varchar(128)        | NO   |     |         |       |
| population | bigint(10) unsigned | YES  |     | NULL    |       |
| area_sq_km | int(9) unsigned     | YES  |     | NULL    |       |
| drive_side | varchar(1)          | YES  |     | NULL    |       |
+------------+---------------------+------+-----+---------+-------+
```

15.19.6.2 Adapting a memcached Application for the InnoDB memcached Plugin

Consider these aspects of MySQL and `InnoDB` tables when adapting existing `memcached` applications to use the `daemon_memcached` plugin:

- If there are key values longer than a few bytes, it may be more efficient to use a numeric auto-increment column as the primary key of the `InnoDB` table, and to create a unique secondary index on the column that contains the `memcached` key values. This is because `InnoDB` performs best for large-scale insertions if primary key values are added in sorted order (as they are with auto-increment values). Primary key values are included in secondary indexes, which takes up unnecessary space if the primary key is a long string value.
- If you store several different classes of information using `memcached`, consider setting up a separate `InnoDB` table for each type of data. Define additional table identifiers in the `innodb_memcache.containers` table, and use the `@@table_id.key` notation to store and retrieve items from different tables. Physically dividing different types of information allows you tune the characteristics of each table for optimum space utilization, performance, and reliability. For example, you might enable compression for a table that holds blog posts, but not for a table that holds thumbnail images. You might back up one table more frequently than another because it holds critical data. You might create additional secondary indexes on tables that are frequently used to generate reports using SQL.
- Preferably, configure a stable set of table definitions for use with the `daemon_memcached` plugin, and leave the tables in place permanently. Changes to the `innodb_memcache.containers` table take effect the next time the `innodb_memcache.containers` table is queried. Entries in the containers table are processed at startup, and are consulted whenever an unrecognized table identifier (as defined by `containers.name`) is requested using `@@` notation. Thus, new entries are visible as soon as you use the associated table identifier, but changes to existing entries require a server restart before they take effect.
- When you use the default `innodb_only` caching policy, calls to `add()`, `set()`, `incr()`, and so on can succeed but still trigger debugging messages such as `while expecting 'STORED', got unexpected response 'NOT_STORED`. Debug messages occur because new and updated values are sent directly to the `InnoDB` table without being saved in the memory cache, due to the `innodb_only` caching policy.

15.19.6.3 Tuning InnoDB memcached Plugin Performance

Because using `InnoDB` in combination with `memcached` involves writing all data to disk, whether immediately or sometime later, raw performance is expected to be somewhat slower than using `memcached` by itself. When using the `InnoDB` `memcached` plugin, focus tuning goals for `memcached` operations on achieving better performance than equivalent SQL operations.

Benchmarks suggest that queries and DML operations (inserts, updates, and deletes) that use the `memcached` interface are faster than traditional SQL. DML operations typically see a larger improvements. Therefore, consider adapting write-intensive applications to use the `memcached` interface first. Also consider prioritizing adaptation of write-intensive applications that use fast, lightweight mechanisms that lack reliability.

Adapting SQL Queries

The types of queries that are most suited to simple `GET` requests are those with a single clause or a set of `AND` conditions in the `WHERE` clause:

```
SQL:
SELECT col FROM tbl WHERE key = 'key_value';

memcached:
GET key_value

SQL:
SELECT col FROM tbl WHERE col1 = val1 and col2 = val2 and col3 = val3;

memcached:
# Since you must always know these 3 values to look up the key,
# combine them into a unique string and use that as the key
# for all ADD, SET, and GET operations.
key_value = val1 + ":" + val2 + ":" + val3
GET key_value

SQL:
SELECT 'key exists!' FROM tbl
  WHERE EXISTS (SELECT col1 FROM tbl WHERE KEY = 'key_value') LIMIT 1;

memcached:
# Test for existence of key by asking for its value and checking if the call succeeds,
# ignoring the value itself. For existence checking, you typically only store a very
# short value such as "1".
GET key_value
```

Using System Memory

For best performance, deploy the `daemon_memcached` plugin on machines that are configured as typical database servers, where the majority of system RAM is devoted to the `InnoDB` buffer pool, through the `innodb_buffer_pool_size` configuration option. For systems with multi-gigabyte buffer pools, consider raising the value of `innodb_buffer_pool_instances` for maximum throughput when most operations involve data that is already cached in memory.

Reducing Redundant I/O

`InnoDB` has a number of settings that let you choose the balance between high reliability, in case of a crash, and the amount of I/O overhead during high write workloads. For example, consider setting the `innodb_doublewrite` to `0` and `innodb_flush_log_at_trx_commit` to `2`. Measure performance with different `innodb_flush_method` settings.

For other ways to reduce or tune I/O for table operations, see Section 8.5.8, “Optimizing InnoDB Disk I/O”.

Reducing Transactional Overhead

A default value of 1 for `daemon_memcached_r_batch_size` and `daemon_memcached_w_batch_size` is intended for maximum reliability of results and safety of stored or updated data.

Depending on the type of application, you might increase one or both of these settings to reduce the overhead of frequent commit operations. On a busy system, you might increase `daemon_memcached_r_batch_size`, knowing that changes to data made through SQL may not become visible to `memcached` immediately (that is, until *`N`* more `get` operations are processed). When processing data where every write operation must be reliably stored, leave `daemon_memcached_w_batch_size` set to `1`. Increase the setting when processing large numbers of updates intended only for statistical analysis, where losing the last *`N`* updates in a crash is an acceptable risk.

For example, imagine a system that monitors traffic crossing a busy bridge, recording data for approximately 100,000 vehicles each day. If the application counts different types of vehicles to analyze traffic patterns, changing `daemon_memcached_w_batch_size` from `1` to `100` reduces I/O overhead for commit operations by 99%. In case of an outage, a maximum of 100 records are

lost, which may be an acceptable margin of error. If instead the application performed automated toll collection for each car, you would set `daemon_memcached_w_batch_size` to `1` to ensure that each toll record is immediately saved to disk.

Because of the way `InnoDB` organizes `memcached` key values on disk, if you have a large number of keys to create, it may be faster to sort the data items by key value in the application and `add` them in sorted order, rather than create keys in arbitrary order.

The `memslap` command, which is part of the regular `memcached` distribution but not included with the `daemon_memcached` plugin, can be useful for benchmarking different configurations. It can also be used to generate sample key/value pairs to use in your own benchmarks. See `libmemcached` Command-Line Utilities for details.

15.19.6.4 Controlling Transactional Behavior of the InnoDB memcached Plugin

Unlike traditional `memcached`, the `daemon_memcached` plugin allows you to control durability of data values produced through calls to `add`, `set`, `incr`, and so on. By default, data written through the `memcached` interface is stored to disk, and calls to `get` return the most recent value from disk. Although the default behavior does not offer the best possible raw performance, it is still fast compared to the SQL interface for `InnoDB` tables.

As you gain experience using the `daemon_memcached` plugin, you can consider relaxing durability settings for non-critical classes of data, at the risk of losing some updated values in the event of an outage, or returning data that is slightly out-of-date.

Frequency of Commits

One tradeoff between durability and raw performance is how frequently new and changed data is committed. If data is critical, is should be committed immediately so that it is safe in case of a crash or outage. If data is less critical, such as counters that are reset after a crash or logging data that you can afford to lose, you might prefer higher raw throughput that is available with less frequent commits.

When a `memcached` operation inserts, updates, or deletes data in the underlying `InnoDB` table, the change might be committed to the `InnoDB` table instantly (if `daemon_memcached_w_batch_size=1`) or some time later (if the `daemon_memcached_w_batch_size` value is greater than 1). In either case, the change cannot be rolled back. If you increase the value of `daemon_memcached_w_batch_size` to avoid high I/O overhead during busy times, commits could become infrequent when the workload decreases. As a safety measure, a background thread automatically commits changes made through the `memcached` API at regular intervals. The interval is controlled by the `innodb_api_bk_commit_interval` configuration option, which has a default setting of `5` seconds.

When a `memcached` operation inserts or updates data in the underlying `InnoDB` table, the changed data is immediately visible to other `memcached` requests because the new value remains in the memory cache, even if it is not yet committed on the MySQL side.

Transaction Isolation

When a `memcached` operation such as `get` or `incr` causes a query or DML operation on the underlying `InnoDB` table, you can control whether the operation sees the very latest data written to the table, only data that has been committed, or other variations of transaction isolation level. Use the `innodb_api_trx_level` configuration option to control this feature. The numeric values specified for this option correspond to isolation levels such as `REPEATABLE READ`. See the description of the `innodb_api_trx_level` option for information about other settings.

A strict isolation level ensures that data you retrieve is not rolled back or changed suddenly causing subsequent queries to return different values. However, strict isolation levels require greater locking overhead, which can cause waits. For a NoSQL-style application that does not use long-running transactions, you can typically use the default isolation level or switch to a less strict isolation level.

Disabling Row Locks for memcached DML Operations

The `innodb_api_disable_rowlock` option can be used to disable row locks when `memcached` requests through the `daemon_memcached` plugin cause DML operations. By default, `innodb_api_disable_rowlock` is set to `OFF` which means that `memcached` requests row locks for `get` and `set` operations. When `innodb_api_disable_rowlock` is set to `ON`, `memcached` requests a table lock instead of row locks.

The `innodb_api_disable_rowlock` option is not dynamic. It must be specified at startup on the `mysqld` command line or entered in a MySQL configuration file.

Allowing or Disallowing DDL

By default, you can perform DDL operations such as `ALTER TABLE` on tables used by the `daemon_memcached` plugin. To avoid potential slowdowns when these tables are used for high-throughput applications, disable DDL operations on these tables by enabling `innodb_api_enable_mdl` at startup. This option is less appropriate when accessing the same tables through both `memcached` and SQL, because it blocks `CREATE INDEX` statements on the tables, which could be important for running reporting queries.

Storing Data on Disk, in Memory, or Both

The `innodb_memcache.cache_policies` table specifies whether to store data written through the `memcached` interface to disk (`innodb_only`, the default); in memory only, as with traditional `memcached` (`cache-only`); or both (`caching`).

With the `caching` setting, if `memcached` cannot find a key in memory, it searches for the value in an `InnoDB` table. Values returned from `get` calls under the `caching` setting could be out-of-date if the values were updated on disk in the `InnoDB` table but are not yet expired from the memory cache.

The caching policy can be set independently for `get`, `set` (including `incr` and `decr`), `delete`, and `flush` operations.

For example, you might allow `get` and `set` operations to query or update a table and the `memcached` memory cache at the same time (using the `caching` setting), while making `delete`, `flush`, or both operate only on the in-memory copy (using the `cache_only` setting). That way, deleting or flushing an item only expires the item from the cache, and the latest value is returned from the `InnoDB` table the next time the item is requested.

```
mysql> SELECT * FROM innodb_memcache.cache_policies;
+--------------+-------------+-------------+---------------+--------------+
| policy_name  | get_policy  | set_policy  | delete_policy | flush_policy |
+--------------+-------------+-------------+---------------+--------------+
| cache_policy | innodb_only | innodb_only | innodb_only   | innodb_only  |
+--------------+-------------+-------------+---------------+--------------+

mysql> UPDATE innodb_memcache.cache_policies SET set_policy = 'caching'
       WHERE policy_name = 'cache_policy';
```

`innodb_memcache.cache_policies` values are only read at startup. After changing values in this table, uninstall and reinstall the `daemon_memcached` plugin to ensure that changes take effect.

```
mysql> UNINSTALL PLUGIN daemon_memcached;

mysql> INSTALL PLUGIN daemon_memcached soname "libmemcached.so";
```

15.19.6.5 Adapting DML Statements to memcached Operations

Benchmarks suggest that the `daemon_memcached` plugin speeds up DML operations (inserts, updates, and deletes) more than it speeds up queries. Therefore, consider focussing initial development efforts on write-intensive applications that are I/O-bound, and look for opportunities to use MySQL with the `daemon_memcached` plugin for new write-intensive applications.

Single-row DML statements are the easiest types of statements to turn into `memcached` operations. `INSERT` becomes `add`, `UPDATE` becomes `set`, `incr` or `decr`, and `DELETE` becomes `delete`. These operations are guaranteed to only affect one row when issued through the `memcached` interface, because the *`key`* is unique within the table.

In the following SQL examples, `t1` refers to the table used for `memcached` operations, based on the configuration in the `innodb_memcache.containers` table. `key` refers to the column listed under `key_columns`, and `val` refers to the column listed under `value_columns`.

```
INSERT INTO t1 (key,val) VALUES (some_key,some_value);
SELECT val FROM t1 WHERE key = some_key;
UPDATE t1 SET val = new_value WHERE key = some_key;
UPDATE t1 SET val = val + x WHERE key = some_key;
DELETE FROM t1 WHERE key = some_key;
```

The following `TRUNCATE TABLE` and `DELETE` statements, which remove all rows from the table, correspond to the `flush_all` operation, where `t1` is configured as the table for `memcached` operations, as in the previous example.

```
TRUNCATE TABLE t1;
DELETE FROM t1;
```

15.19.6.6 Performing DML and DDL Statements on the Underlying InnoDB Table

You can access the underlying `InnoDB` table (which is `test.demo_test` by default) through standard SQL interfaces. However, there are some restrictions:

- When querying a table that is also accessed through the `memcached` interface, remember that `memcached` operations can be configured to be committed periodically rather than after every write operation. This behavior is controlled by the `daemon_memcached_w_batch_size` option. If this option is set to a value greater than `1`, use `READ UNCOMMITTED` queries to find rows that were just inserted.

```
mysql> SET SESSSION TRANSACTION ISOLATION LEVEL READ UNCOMMITTED;

mysql> SELECT * FROM demo_test;
+------+------+------+------+-----------+------+------+------+------+------+------+
| cx   | cy   | c1   | cz   | c2        | ca   | CB   | c3   | cu   | c4   | C5   |
+------+------+------+------+-----------+------+------+------+------+------+------+
| NULL | NULL | a11  | NULL | 123456789 | NULL | NULL |   10 | NULL |    3 | NULL |
+------+------+------+------+-----------+------+------+------+------+------+------+
```

- When modifying a table using SQL that is also accessed through the `memcached` interface, you can configure `memcached` operations to start a new transaction periodically rather than for every read operation. This behavior is controlled by the `daemon_memcached_r_batch_size` option. If this option is set to a value greater than `1`, changes made to the table using SQL are not immediately visible to `memcached` operations.

- The `InnoDB` table is either IS (intention shared) or IX (intention exclusive) locked for all operations in a transaction. If you increase `daemon_memcached_r_batch_size` and `daemon_memcached_w_batch_size` substantially from their default value of `1`, the table is most likely locked between each operation, preventing DDL statements on the table.

15.19.7 The InnoDB memcached Plugin and Replication

Because the `daemon_memcached` plugin supports the MySQL binary log, updates made on a master server through the `memcached` interface can be replicated for backup, balancing intensive read workloads, and high availability. All `memcached` commands are supported with binary logging.

You do not need to set up the `daemon_memcached` plugin on slave servers. The primary advantage of this configuration is increased write throughput on the master. The speed of the replication mechanism is not affected.

The following sections show how to use the binary log capability when using the `daemon_memcached` plugin with MySQL replication. It is assumed that you have completed the setup described in Section 15.19.3, “Setting Up the InnoDB memcached Plugin”.

Enabling the InnoDB memcached Binary Log

1. To use the `daemon_memcached` plugin with the MySQL binary log, enable the `innodb_api_enable_binlog` configuration option on the master server. This option can only be set at server startup. You must also enable the MySQL binary log on the master server using the `--log-bin` option. You can add these options to the MySQL configuration file, or on the `mysqld` command line.

   ```
   mysqld ... --log-bin --innodb_api_enable_binlog=1
   ```

2. Configure the master and slave server, as described in Section 18.1.2, “Setting Up Binary Log File Position Based Replication”.

3. Use `mysqldump` to create a master data snapshot, and sync the snapshot to the slave server.

   ```
   master shell> mysqldump --all-databases --lock-all-tables > dbdump.db
   slave shell> mysql < dbdump.db
   ```

4. On the master server, issue `SHOW MASTER STATUS` to obtain the master binary log coordinates.

   ```
   mysql> SHOW MASTER STATUS;
   ```

5. On the slave server, use a `CHANGE MASTER TO` statement to set up a slave server using the master binary log coordinates.

   ```
   mysql> CHANGE MASTER TO
          MASTER_HOST='localhost',
          MASTER_USER='root',
          MASTER_PASSWORD='',
          MASTER_PORT = 13000,
          MASTER_LOG_FILE='0.000001,
          MASTER_LOG_POS=114;
   ```

6. Start the slave.

   ```
   mysql> START SLAVE;
   ```

   If the error log prints output similar to the following, the slave is ready for replication.

   ```
   2013-09-24T13:04:38.639684Z 49 [Note] Slave I/O thread: connected to
   master 'root@localhost:13000', replication started in log '0.000001'
   at position 114
   ```

Testing the InnoDB memcached Replication Configuration

This example demonstrates how to test the `InnoDB memcached` replication configuration using the `memcached` and telnet to insert, update, and delete data. A MySQL client is used to verify results on the master and slave servers.

The example uses the `demo_test` table, which was created by the `innodb_memcached_config.sql` configuration script during the initial setup of the `daemon_memcached` plugin. The `demo_test` table contains a single example record.

1. Use the `set` command to insert a record with a key of `test1`, a flag value of `10`, an expiration value of `0`, a cas value of 1, and a value of `t1`.

   ```
   ```

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
set test1 10 0 1
t1
STORED
```

2. On the master server, check that the record was inserted into the `demo_test` table. Assuming the `demo_test` table was not previously modified, there should be two records. The example record with a key of `AA`, and the record you just inserted, with a key of `test1`. The `c1` column maps to the key, the `c2` column to the value, the `c3` column to the flag value, the `c4` column to the cas value, and the `c5` column to the expiration time. The expiration time was set to 0, since it is unused.

```
mysql> SELECT * FROM test.demo_test;
+-------+--------------+------+------+------+
| c1    | c2           | c3   | c4   | c5   |
+-------+--------------+------+------+------+
| AA    | HELLO, HELLO |    8 |    0 |    0 |
| test1 | t1           |   10 |    1 |    0 |
+-------+--------------+------+------+------+
```

3. Check to verify that the same record was replicated to the slave server.

```
mysql> SELECT * FROM test.demo_test;
+-------+--------------+------+------+------+
| c1    | c2           | c3   | c4   | c5   |
+-------+--------------+------+------+------+
| AA    | HELLO, HELLO |    8 |    0 |    0 |
| test1 | t1           |   10 |    1 |    0 |
+-------+--------------+------+------+------+
```

4. Use the `set` command to update the key to a value of `new`.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
set test1 10 0 2
new
STORED
```

The update is replicated to the slave server (notice that the `cas` value is also updated).

```
mysql> SELECT * FROM test.demo_test;
+-------+--------------+------+------+------+
| c1    | c2           | c3   | c4   | c5   |
+-------+--------------+------+------+------+
| AA    | HELLO, HELLO |    8 |    0 |    0 |
| test1 | new          |   10 |    2 |    0 |
+-------+--------------+------+------+------+
```

5. Delete the `test1` record using a `delete` command.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
delete test1
DELETED
```

When the `delete` operation is replicated to the slave, the `test1` record on the slave is also deleted.

```
mysql> SELECT * FROM test.demo_test;
+----+--------------+------+------+------+
| c1 | c2           | c3   | c4   | c5   |
+----+--------------+------+------+------+
```

```
| AA | HELLO, HELLO |    8 |    0 |    0 |
+----+--------------+------+------+------+
```

6. Remove all rows from the table using the `flush_all` command.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
flush_all
OK
```

```
mysql> SELECT * FROM test.demo_test;
Empty set (0.00 sec)
```

7. Telnet to the master server and enter two new records.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'
set test2 10 0 4
again
STORED
set test3 10 0 5
again1
STORED
```

8. Confirm that the two records were replicated to the slave server.

```
mysql> SELECT * FROM test.demo_test;
+-------+--------------+------+------+------+
| c1    | c2           | c3   | c4   | c5   |
+-------+--------------+------+------+------+
| test2 | again        |   10 |    4 |    0 |
| test3 | again1       |   10 |    5 |    0 |
+-------+--------------+------+------+------+
```

9. Remove all rows from the table using the `flush_all` command.

```
telnet 127.0.0.1 11211
Trying 127.0.0.1...
Connected to 127.0.0.1.
Escape character is '^]'.
flush_all
OK
```

10. Check to ensure that the `flush_all` operation was replicated on the slave server.

```
mysql> SELECT * FROM test.demo_test;
Empty set (0.00 sec)
```

InnoDB memcached Binary Log Notes

Binary Log Format:

- Most `memcached` operations are mapped to DML statements (analogous to insert, delete, update). Since there is no actual SQL statement being processed by the MySQL server, all `memcached` commands (except for `flush_all`) use Row-Based Replication (RBR) logging, which is independent of any server `binlog_format` setting.
- The `memcached` `flush_all` command is mapped to the `TRUNCATE TABLE` command in MySQL 5.7 and earlier. Since DDL commands can only use statement-based logging, the `flush_all` command is replicated by sending a `TRUNCATE TABLE` statement. In MySQL 8.0 and later, `flush_all` is mapped to `DELETE` but is still replicated by sending a `TRUNCATE TABLE` statement.

Transactions:

- The concept of transactions has not typically been part of `memcached` applications. For performance considerations, `daemon_memcached_r_batch_size` and `daemon_memcached_w_batch_size` are used to control the batch size for read and write transactions. These settings do not affect replication. Each SQL operation on the underlying `InnoDB` table is replicated after successful completion.
- The default value of `daemon_memcached_w_batch_size` is `1`, which means that each `memcached` write operation is committed immediately. This default setting incurs a certain amount of performance overhead to avoid inconsistencies in the data that is visible on the master and slave servers. The replicated records are always available immediately on the slave server. If you set `daemon_memcached_w_batch_size` to a value greater than `1`, records inserted or updated through `memcached` are not immediately visible on the master server; to view the records on the master server before they are committed, issue `SET TRANSACTION ISOLATION LEVEL READ UNCOMMITTED`.

15.19.8 InnoDB memcached Plugin Internals

InnoDB API for the InnoDB memcached Plugin

The `InnoDB` `memcached` engine accesses `InnoDB` through `InnoDB` APIs, most of which are directly adopted from embedded `InnoDB`. `InnoDB` API functions are passed to the `InnoDB` `memcached` engine as callback functions. `InnoDB` API functions access the `InnoDB` tables directly, and are mostly DML operations with the exception of `TRUNCATE TABLE`.

`memcached` commands are implemented through the `InnoDB` `memcached` API. The following table outlines how `memcached` commands are mapped to DML or DDL operations.

Table 15.15 memcached Commands and Associated DML or DDL Operations

memcached Command	DML or DDL Operations
`get`	a read/fetch command
`set`	a search followed by an `INSERT` or `UPDATE` (depending on whether or not a key exists)
`add`	a search followed by an `INSERT` or `UPDATE`
`replace`	a search followed by an `UPDATE`
`append`	a search followed by an `UPDATE` (appends data to the result before `UPDATE`)
`prepend`	a search followed by an `UPDATE` (prepends data to the result before `UPDATE`)
`incr`	a search followed by an `UPDATE`
`decr`	a search followed by an `UPDATE`
`delete`	a search followed by a `DELETE`
`flush_all`	`TRUNCATE TABLE` (DDL)

InnoDB memcached Plugin Configuration Tables

This section describes configuration tables used by the `daemon_memcached` plugin. The `cache_policies` table, `config_options` table, and `containers` table are created by the `innodb_memcached_config.sql` configuration script in the `innodb_memcache` database.

```
mysql> USE innodb_memcache;
Database changed
mysql> SHOW TABLES;
+---------------------------+
| Tables_in_innodb_memcache |
+---------------------------+
```

```
| cache_policies            |
| config_options            |
| containers                |
+---------------------------+
```

cache_policies Table

The `cache_policies` table defines a cache policy for the `InnoDB` `memcached` installation. You can specify individual policies for `get`, `set`, `delete`, and `flush` operations, within a single cache policy. The default setting for all operations is `innodb_only`.

- `innodb_only`: Use `InnoDB` as the data store.
- `cache-only`: Use the `memcached` engine as the data store.
- `caching`: Use both `InnoDB` and the `memcached` engine as data stores. In this case, if `memcached` cannot find a key in memory, it searches for the value in an `InnoDB` table.
- `disable`: Disable caching.

Table 15.16 cache_policies Columns

Column	Description
`policy_name`	Name of the cache policy. The default cache policy name is `cache_policy`.
`get_policy`	The cache policy for get operations. Valid values are `innodb_only`, `cache-only`, `caching`, or `disabled`. The default setting is `innodb_only`.
`set_policy`	The cache policy for set operations. Valid values are `innodb_only`, `cache-only`, `caching`, or `disabled`. The default setting is `innodb_only`.
`delete_policy`	The cache policy for delete operations. Valid values are `innodb_only`, `cache-only`, `caching`, or `disabled`. The default setting is `innodb_only`.
`flush_policy`	The cache policy for flush operations. Valid values are `innodb_only`, `cache-only`, `caching`, or `disabled`. The default setting is `innodb_only`.

config_options Table

The `config_options` table stores `memcached`-related settings that can be changed at runtime using SQL. Supported configuration options are `separator` and `table_map_delimiter`.

Table 15.17 config_options Columns

Column	Description
`Name`	Name of the `memcached`-related configuration option. The following configuration options are supported by the `config_options` table: • `separator`: Used to separate values of a long string into separate values when there are multiple `value_columns` defined. By default, the `separator` is a `|` character. For example, if you define `col1, col2` as value columns, and you define `|` as the separator, you can issue the following `memcached` command to insert values into `col1` and `col2`, respectively:

Column	Description

	`set keyx 10 0 19`
	`valuecolx|valuecoly`
	
	`valuecol1x` is stored in `col1` and `valuecoly` is stored in `col2`.
	
	• `table_map_delimiter`: The character separating the schema name and the table name when you use the `@@` notation in a key name to access a key in a specific table. For example, `@@t1.some_key` and `@@t2.some_key` have the same key value, but are stored in different tables.

`Value`	The value assigned to the `memcached`-related configuration option.

containers Table

The `containers` table is the most important of the three configuration tables. Each `InnoDB` table that is used to store `memcached` values must have an entry in the `containers` table. The entry provides a mapping between `InnoDB` table columns and container table columns, which is required for `memcached` to work with `InnoDB` tables.

The `containers` table contains a default entry for the `test.demo_test` table, which is created by the `innodb_memcached_config.sql` configuration script. To use the `daemon_memcached` plugin with your own `InnoDB` table, you must create an entry in the `containers` table.

Table 15.18 containers Columns

Column	Description
`name`	The name given to the container. If an `InnoDB` table is not requested by name using `@@` notation, the `daemon_memcached` plugin uses the `InnoDB` table with a `containers.name` value of `default`. If there is no such entry, the first entry in the `containers` table, ordered alphabetically by `name` (ascending), determines the default `InnoDB` table.
`db_schema`	The name of the database where the `InnoDB` table resides. This is a required value.
`db_table`	The name of the `InnoDB` table that stores `memcached` values. This is a required value.
`key_columns`	The column in the `InnoDB` table that contains lookup key values for `memcached` operations. This is a required value.
`value_columns`	The `InnoDB` table columns (one or more) that store `memcached` data. Multiple columns can be specified using the separator character specified in the `innodb_memcached.config_options` table. By default, the separator is a pipe character ("|"). To specify multiple columns, separate them with the defined separator character. For example: `col1|col2|col3`. This is a required value.
`flags`	The `InnoDB` table columns that are used as flags (a user-defined numeric value that is stored

Column	Description
	and retrieved along with the main value) for `memcached`. A flag value can be used as a column specifier for some operations (such as `incr`, `prepend`) if a `memcached` value is mapped to multiple columns, so that an operation is performed on a specified column. For example, if you have mapped a `value_columns` to three `InnoDB` table columns, and only want the increment operation performed on one columns, use the `flags` column to specify the column. If you do not use the `flags` column, set a value of `0` to indicate that it is unused.
`cas_column`	The `InnoDB` table column that stores compare-and-swap (cas) values. The `cas_column` value is related to the way `memcached` hashes requests to different servers and caches data in memory. Because the `InnoDB` `memcached` plugin is tightly integrated with a single `memcached` daemon, and the in-memory caching mechanism is handled by MySQL and the InnoDB buffer pool, this column is rarely needed. If you do not use this column, set a value of `0` to indicate that it is unused.
`expire_time_column`	The `InnoDB` table column that stores expiration values. The `expire_time_column` value is related to the way `memcached` hashes requests to different servers and caches data in memory. Because the `InnoDB` `memcached` plugin is tightly integrated with a single `memcached` daemon, and the in-memory caching mechanism is handled by MySQL and the InnoDB buffer pool, this column is rarely needed. If you do not use this column, set a value of `0` to indicate that the column is unused. The maximum expire time is defined as `INT_MAX32` or 2147483647 seconds (approximately 68 years).
`unique_idx_name_on_key`	The name of the index on the key column. It must be a unique index. It can be the primary key or a secondary index. Preferably, use the primary key of the `InnoDB` table. Using the primary key avoids a lookup that is performed when using a secondary index. You cannot make a covering index for `memcached` lookups; `InnoDB` returns an error if you try to define a composite secondary index over both the key and value columns.

containers Table Column Constraints

- You must supply a value for `db_schema`, `db_name`, `key_columns`, `value_columns` and `unique_idx_name_on_key`. Specify `0` for `flags`, `cas_column`, and `expire_time_column` if they are unused. Failing to do so could cause your setup to fail.
- `key_columns`: The maximum limit for a `memcached` key is 250 characters, which is enforced by `memcached`. The mapped key must be a non-Null `CHAR` or `VARCHAR` type.
- `value_columns`: Must be mapped to a `CHAR`, `VARCHAR`, or `BLOB` column. There is no length restriction and the value can be NULL.

- `cas_column`: The `cas` value is a 64 bit integer. It must be mapped to a `BIGINT` of at least 8 bytes. If you do not use this column, set a value of `0` to indicate that it is unused.
- `expiration_time_column`: Must mapped to an `INTEGER` of at least 4 bytes. Expiration time is defined as a 32-bit integer for Unix time (the number of seconds since January 1, 1970, as a 32-bit value), or the number of seconds starting from the current time. For the latter, the number of seconds may not exceed 60*60*24*30 (the number of seconds in 30 days). If the number sent by a client is larger, the server considers it to be a real Unix time value rather than an offset from the current time. If you do not use this column, set a value of `0` to indicate that it is unused.
- `flags`: Must be mapped to an `INTEGER` of at least 32-bits and can be NULL. If you do not use this column, set a value of `0` to indicate that it is unused.

A pre-check is performed at plugin load time to enforce column constraints. If mismatches are found, the plugin is not loaded.

Multiple Value Column Mapping

- During plugin initialization, when `InnoDB memcached` is configured with information defined in the `containers` table, each mapped column defined in `containers.value_columns` is verified against the mapped `InnoDB` table. If multiple `InnoDB` table columns are mapped, there is a check to ensure that each column exists and is the right type.
- At run-time, for `memcached` insert operations, if there are more delimited values than the number of mapped columns, only the number of mapped values are taken. For example, if there are six mapped columns, and seven delimited values are provided, only the first six delimited values are taken. The seventh delimited value is ignored.
- If there are fewer delimited values than mapped columns, unfilled columns are set to NULL. If an unfilled column cannot be set to NULL, insert operations fail.
- If a table has more columns than mapped values, the extra columns do not affect results.

The demo_test Example Table

The `innodb_memcached_config.sql` configuration script creates a `demo_test` table in the `test` database, which can be used to verify `InnoDB memcached` plugin installation immediately after setup.

The `innodb_memcached_config.sql` configuration script also creates an entry for the `demo_test` table in the `innodb_memcache.containers` table.

```
mysql> SELECT * FROM innodb_memcache.containers\G
*************************** 1. row ***************************
                name: aaa
           db_schema: test
            db_table: demo_test
         key_columns: c1
       value_columns: c2
               flags: c3
          cas_column: c4
   expire_time_column: c5
unique_idx_name_on_key: PRIMARY

mysql> SELECT * FROM test.demo_test;
+----+-----------------+------+------+------+
| c1 | c2              | c3   | c4   | c5   |
+----+-----------------+------+------+------+
| AA | HELLO, HELLO    |    8 |    0 |    0 |
+----+-----------------+------+------+------+
```

15.19.9 Troubleshooting the InnoDB memcached Plugin

This section describes issues that you may encounter when using the `InnoDB memcached` plugin.

- If you encounter the following error in the MySQL error log, the server might fail to start:

  ```
  failed to set rlimit for open files. Try running as root or requesting
  smaller maxconns value.
  ```

  The error message is from the `memcached` daemon. One solution is to raise the OS limit for the number of open files. The commands for checking and increasing the open file limit varies by operating system. This example shows commands for Linux and OS X:

  ```
  # Linux
  shell> ulimit -n
  1024
  shell> ulimit -n 4096
  shell> ulimit -n
  4096

  # OS X
  shell> ulimit -n
  256
  shell> ulimit -n 4096
  shell> ulimit -n
  4096
  ```

  The other solution is to reduce the number of concurrent connections permitted for the `memcached` daemon. To do so, encode the `-c memcached` option in the `daemon_memcached_option` configuration parameter in the MySQL configuration file. The `-c` option has a default value of 1024.

  ```
  [mysqld]
  ...
  loose-daemon_memcached_option='-c 64'
  ```

- To troubleshoot problems where the `memcached` daemon is unable to store or retrieve `InnoDB` table data, encode the `-vvv memcached` option in the `daemon_memcached_option` configuration parameter in the MySQL configuration file. Examine the MySQL error log for debug output related to `memcached` operations.

  ```
  [mysqld]
  ...
  loose-daemon_memcached_option='-vvv'
  ```

- If columns specified to hold `memcached` values are the wrong data type, such as a numeric type instead of a string type, attempts to store key/value pairs fail with no specific error code or message.

- If the `daemon_memcached` plugin causes MySQL server startup issues, you can temporarily disable the `daemon_memcached` plugin while troubleshooting by adding this line under the `[mysqld]` group in the MySQL configuration file:

  ```
  daemon_memcached=OFF
  ```

  For example, if you run the `INSTALL PLUGIN` statement before running the `innodb_memcached_config.sql` configuration script to set up the necessary database and tables, the server might crash and fail to start. The server could also fail to start if you incorrectly configure an entry in the `innodb_memcache.containers` table.

  To uninstall the `memcached` plugin for a MySQL instance, issue the following statement:

  ```
  mysql> UNINSTALL PLUGIN daemon_memcached;
  ```

- If you run more than one instance of MySQL on the same machine with the `daemon_memcached` plugin enabled in each instance, use the `daemon_memcached_option` configuration parameter to specify a unique `memcached` port for each `daemon_memcached` plugin.

- If an SQL statement cannot find the `InnoDB` table or finds no data in the table, but `memcached` API calls retrieve the expected data, you may be missing an entry for the `InnoDB` table in the `innodb_memcache.containers` table, or you may have not switched to the correct `InnoDB` table by issuing a `get` or `set` request using `@@table_id` notation. This problem could also occur if you change an existing entry in the `innodb_memcache.containers` table without restarting the MySQL server afterward. The free-form storage mechanism is flexible enough that your requests to store or retrieve a multi-column value such as `col1|col2|col3` may still work, even if the daemon is using the `test.demo_test` table which stores values in a single column.

- When defining your own `InnoDB` table for use with the `daemon_memcached` plugin, and columns in the table are defined as `NOT NULL`, ensure that values are supplied for the `NOT NULL` columns when inserting a record for the table into the `innodb_memcache.containers` table. If the `INSERT` statement for the `innodb_memcache.containers` record contains fewer delimited values than there are mapped columns, unfilled columns are set to `NULL`. Attempting to insert a `NULL` value into a `NOT NULL` column causes the `INSERT` to fail, which may only become evident after you reinitialize the `daemon_memcached` plugin to apply changes to the `innodb_memcache.containers` table.

- If `cas_column` and `expire_time_column` fields of the `innodb_memcached.containers` table are set to `NULL`, the following error is returned when attempting to load the `memcached` plugin:

  ```
  InnoDB_Memcached: column 6 in the entry for config table 'containers' in
  database 'innodb_memcache' has an invalid NULL value.
  ```

  The `memcached` plugin rejects usage of `NULL` in the `cas_column` and `expire_time_column` columns. Set the value of these columns to `0` when the columns are unused.

- As the length of the `memcached` key and values increase, you might encounter size and length limits.

  - When the key exceeds 250 bytes, `memcached` operations return an error. This is currently a fixed limit within `memcached`.

  - `InnoDB` table limits may be encountered if values exceed 768 bytes in size, 3072 bytes in size, or half of the `innodb_page_size` value. These limits primarily apply if you intend to create an index on a value column to run report-generating queries on that column using SQL. See Section 15.8.1.7, “Limits on InnoDB Tables” for details.

  - The maximum size for the key-value combination is 1 MB.

- If you share configuration files across MySQL servers of different versions, using the latest configuration options for the `daemon_memcached` plugin could cause startup errors on older MySQL versions. To avoid compatibility problems, use the `loose` prefix with option names. For example, use `loose-daemon_memcached_option='-c 64'` instead of `daemon_memcached_option='-c 64'`.

- There is no restriction or check in place to validate character set settings. `memcached` stores and retrieves keys and values in bytes and is therefore not character set sensitive. However, you must ensure that the `memcached` client and the MySQL table use the same character set.

- `memcached` connections are blocked from accessing tables that contain an indexed virtual column. Accessing an indexed virtual column requires a callback to the server, but a `memcached` connection does not have access to the server code.

15.20 InnoDB Troubleshooting

The following general guidelines apply to troubleshooting `InnoDB` problems:

- When an operation fails or you suspect a bug, look at the MySQL server error log (see Section 5.4.2, “The Error Log”). Section B.3, “Server Error Codes and Messages” provides troubleshooting information for some of the common `InnoDB`-specific errors that you may encounter.

- If the failure is related to a deadlock, run with the `innodb_print_all_deadlocks` option enabled so that details about each deadlock are printed to the MySQL server error log. For information about deadlocks, see Section 15.5.5, “Deadlocks in InnoDB”.
- If the issue is related to the `InnoDB` data dictionary, see Section 15.20.3, “Troubleshooting InnoDB Data Dictionary Operations”.
- When troubleshooting, it is usually best to run the MySQL server from the command prompt, rather than through `mysqld_safe` or as a Windows service. You can then see what `mysqld` prints to the console, and so have a better grasp of what is going on. On Windows, start `mysqld` with the `--console` option to direct the output to the console window.
- Enable the `InnoDB` Monitors to obtain information about a problem (see Section 15.16, “InnoDB Monitors”). If the problem is performance-related, or your server appears to be hung, you should enable the standard Monitor to print information about the internal state of `InnoDB`. If the problem is with locks, enable the Lock Monitor. If the problem is with table creation, tablespaces, or data dictionary operations, refer to the InnoDB Information Schema system tables to examine contents of the `InnoDB` internal data dictionary.

  `InnoDB` temporarily enables standard `InnoDB` Monitor output under the following conditions:

  - A long semaphore wait
  - `InnoDB` cannot find free blocks in the buffer pool
  - Over 67% of the buffer pool is occupied by lock heaps or the adaptive hash index
- If you suspect that a table is corrupt, run `CHECK TABLE` on that table.

15.20.1 Troubleshooting InnoDB I/O Problems

The troubleshooting steps for `InnoDB` I/O problems depend on when the problem occurs: during startup of the MySQL server, or during normal operations when a DML or DDL statement fails due to problems at the file system level.

Initialization Problems

If something goes wrong when `InnoDB` attempts to initialize its tablespace or its log files, delete all files created by `InnoDB`: all `ibdata` files and all `ib_logfile` files. If you already created some `InnoDB` tables, also delete any `.ibd` files from the MySQL database directories. Then try the `InnoDB` database creation again. For easiest troubleshooting, start the MySQL server from a command prompt so that you see what is happening.

Runtime Problems

If `InnoDB` prints an operating system error during a file operation, usually the problem has one of the following solutions:

- Make sure the `InnoDB` data file directory and the `InnoDB` log directory exist.
- Make sure `mysqld` has access rights to create files in those directories.
- Make sure `mysqld` can read the proper `my.cnf` or `my.ini` option file, so that it starts with the options that you specified.
- Make sure the disk is not full and you are not exceeding any disk quota.
- Make sure that the names you specify for subdirectories and data files do not clash.
- Doublecheck the syntax of the `innodb_data_home_dir` and `innodb_data_file_path` values. In particular, any `MAX` value in the `innodb_data_file_path` option is a hard limit, and exceeding that limit causes a fatal error.

15.20.2 Forcing InnoDB Recovery

To investigate database page corruption, you might dump your tables from the database with `SELECT ... INTO OUTFILE`. Usually, most of the data obtained in this way is intact. Serious corruption might cause `SELECT * FROM tbl_name` statements or `InnoDB` background operations to crash or assert, or even cause `InnoDB` roll-forward recovery to crash. In such cases, you can use the `innodb_force_recovery` option to force the `InnoDB` storage engine to start up while preventing background operations from running, so that you can dump your tables. For example, you can add the following line to the `[mysqld]` section of your option file before restarting the server:

```
[mysqld]
innodb_force_recovery = 1
```

Warning

Only set `innodb_force_recovery` to a value greater than 0 in an emergency situation, so that you can start `InnoDB` and dump your tables. Before doing so, ensure that you have a backup copy of your database in case you need to recreate it. Values of 4 or greater can permanently corrupt data files. Only use an `innodb_force_recovery` setting of 4 or greater on a production server instance after you have successfully tested the setting on a separate physical copy of your database. When forcing `InnoDB` recovery, you should always start with `innodb_force_recovery=1` and only increase the value incrementally, as necessary.

`innodb_force_recovery` is 0 by default (normal startup without forced recovery). The permissible nonzero values for `innodb_force_recovery` are 1 to 6. A larger value includes the functionality of lesser values. For example, a value of 3 includes all of the functionality of values 1 and 2.

If you are able to dump your tables with an `innodb_force_recovery` value of 3 or less, then you are relatively safe that only some data on corrupt individual pages is lost. A value of 4 or greater is considered dangerous because data files can be permanently corrupted. A value of 6 is considered drastic because database pages are left in an obsolete state, which in turn may introduce more corruption into B-trees and other database structures.

As a safety measure, `InnoDB` prevents `INSERT`, `UPDATE`, or `DELETE` operations when `innodb_force_recovery` is greater than 0. An `innodb_force_recovery` setting of 4 or greater places `InnoDB` in read-only mode.

- `1` (`SRV_FORCE_IGNORE_CORRUPT`)

  Lets the server run even if it detects a corrupt page. Tries to make `SELECT * FROM tbl_name` jump over corrupt index records and pages, which helps in dumping tables.

- `2` (`SRV_FORCE_NO_BACKGROUND`)

  Prevents the master thread and any purge threads from running. If a crash would occur during the purge operation, this recovery value prevents it.

- `3` (`SRV_FORCE_NO_TRX_UNDO`)

  Does not run transaction rollbacks after crash recovery.

- `4` (`SRV_FORCE_NO_IBUF_MERGE`)

  Prevents insert buffer merge operations. If they would cause a crash, does not do them. Does not calculate table statistics. This value can permanently corrupt data files. After using this value, be prepared to drop and recreate all secondary indexes. Sets `InnoDB` to read-only.

- `5` (`SRV_FORCE_NO_UNDO_LOG_SCAN`)

Does not look at undo logs when starting the database: InnoDB treats even incomplete transactions as committed. This value can permanently corrupt data files. Sets InnoDB to read-only.

- 6 (SRV_FORCE_NO_LOG_REDO)

  Does not do the redo log roll-forward in connection with recovery. This value can permanently corrupt data files. Leaves database pages in an obsolete state, which in turn may introduce more corruption into B-trees and other database structures. Sets InnoDB to read-only.

You can SELECT from tables to dump them. With an innodb_force_recovery value of 3 or less you can DROP or CREATE tables. DROP TABLE is also supported with an innodb_force_recovery value greater than 3. DROP TABLE is not permitted with an innodb_force_recovery value greater than 4.

If you know that a given table is causing a crash on rollback, you can drop it. If you encounter a runaway rollback caused by a failing mass import or ALTER TABLE, you can kill the mysqld process and set innodb_force_recovery to 3 to bring the database up without the rollback, and then DROP the table that is causing the runaway rollback.

If corruption within the table data prevents you from dumping the entire table contents, a query with an ORDER BY *primary_key* DESC clause might be able to dump the portion of the table after the corrupted part.

If a high innodb_force_recovery value is required to start InnoDB, there may be corrupted data structures that could cause complex queries (queries containing WHERE, ORDER BY, or other clauses) to fail. In this case, you may only be able to run basic SELECT * FROM t queries.

15.20.3 Troubleshooting InnoDB Data Dictionary Operations

Information about table definitions is stored in the InnoDB data dictionary. If you move data files around, or if the server crashes in the middle of a data dictionary operation, dictionary data can become inconsistent.

If a data dictionary corruption or consistency issue prevents you from starting InnoDB, see Section 15.20.2, “Forcing InnoDB Recovery” for information about manual recovery.

Cannot Open Datafile

With innodb_file_per_table enabled (the default), the following messages may appear at startup if a file-per-table tablespace file (.ibd file) is missing:

```
[ERROR] InnoDB: Operating system error number 2 in a file operation.
[ERROR] InnoDB: The error means the system cannot find the path specified.
[ERROR] InnoDB: Cannot open datafile for read-only: './test/t1.ibd' OS error: 71
[Warning] InnoDB: Ignoring tablespace `test/t1` because it could not be opened.
```

To address the these messages, issue DROP TABLE statement to remove data about the missing table from the data dictionary.

Restoring Orphan File-Per-Table ibd Files

This procedure describes how to restore orphan file-per-table .ibd files to another MySQL instance. You might use this procedure if the system tablespace is lost or unrecoverable and you want to restore .idb file backups on a new MySQL instance.

The procedure is not supported for general tablespace .ibd files.

The procedure assumes that you only have .ibd file backups, you are recovering to the same version of MySQL that initially created the orphan .idb files, and that .idb file backups are clean. See Section 15.8.1.3, “Moving or Copying InnoDB Tables” for information about creating clean backups.

Tablespace copying limitations outlined in Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance” are applicable to this procedure.

1. On the new MySQL instance, recreate the table in a database of the same name.

```
mysql> CREATE DATABASE sakila;

mysql> USE sakila;

mysql> CREATE TABLE actor (
    ->    actor_id SMALLINT UNSIGNED NOT NULL AUTO_INCREMENT,
    ->    first_name VARCHAR(45) NOT NULL,
    ->    last_name VARCHAR(45) NOT NULL,
    ->    last_update TIMESTAMP NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP,
    ->    PRIMARY KEY  (actor_id),
    ->    KEY idx_actor_last_name (last_name)
    -> )ENGINE=InnoDB DEFAULT CHARSET=utf8;
```

2. Discard the tablespace of the newly created table.

```
mysql> ALTER TABLE sakila.actor DISCARD TABLESPACE;
```

3. Copy the orphan `.idb` file from your backup directory to the new database directory.

```
shell> cp /backup_directory/actor.ibd path/to/mysql-5.7/data/sakila/
```

4. Ensure that the `.ibd` file has the necessary file permissions.

5. Import the orphan `.ibd` file. A warning is issued indicating that `InnoDB` will attempt to import the file without schema verification.

```
mysql> ALTER TABLE sakila.actor IMPORT TABLESPACE; SHOW WARNINGS;
Query OK, 0 rows affected, 1 warning (0.15 sec)

Warning | 1810 | InnoDB: IO Read error: (2, No such file or directory)
Error opening './sakila/actor.cfg', will attempt to import
without schema verification
```

6. Query the table to verify that the `.ibd` file was successfully restored.

```
mysql> SELECT COUNT(*) FROM sakila.actor;
+----------+
| count(*) |
+----------+
|      200 |
+----------+
```

15.20.4 InnoDB Error Handling

The following items describe how `InnoDB` performs error handling. `InnoDB` sometimes rolls back only the statement that failed, other times it rolls back the entire transaction.

- If you run out of file space in a tablespace, a MySQL `Table is full` error occurs and `InnoDB` rolls back the SQL statement.

- A transaction deadlock causes `InnoDB` to roll back the entire transaction. Retry the whole transaction when this happens.

  A lock wait timeout causes `InnoDB` to roll back only the single statement that was waiting for the lock and encountered the timeout. (To have the entire transaction roll back, start the server with the `--innodb_rollback_on_timeout` option.) Retry the statement if using the current behavior, or the entire transaction if using `--innodb_rollback_on_timeout`.

  Both deadlocks and lock wait timeouts are normal on busy servers and it is necessary for applications to be aware that they may happen and handle them by retrying. You can make them less likely by doing as little work as possible between the first change to data during a transaction

and the commit, so the locks are held for the shortest possible time and for the smallest possible number of rows. Sometimes splitting work between different transactions may be practical and helpful.

When a transaction rollback occurs due to a deadlock or lock wait timeout, it cancels the effect of the statements within the transaction. But if the start-transaction statement was `START TRANSACTION` or `BEGIN` statement, rollback does not cancel that statement. Further SQL statements become part of the transaction until the occurrence of `COMMIT`, `ROLLBACK`, or some SQL statement that causes an implicit commit.

- A duplicate-key error rolls back the SQL statement, if you have not specified the `IGNORE` option in your statement.
- A `row too long error` rolls back the SQL statement.
- Other errors are mostly detected by the MySQL layer of code (above the `InnoDB` storage engine level), and they roll back the corresponding SQL statement. Locks are not released in a rollback of a single SQL statement.

During implicit rollbacks, as well as during the execution of an explicit `ROLLBACK` SQL statement, `SHOW PROCESSLIST` displays `Rolling back` in the `State` column for the relevant connection.

Chapter 16 Alternative Storage Engines

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Storage engines are MySQL components that handle the SQL operations for different table types. `InnoDB` is the default and most general-purpose storage engine, and Oracle recommends using it for tables except for specialized use cases. (The `CREATE TABLE` statement in MySQL 8.0 creates `InnoDB` tables by default.)

MySQL Server uses a pluggable storage engine architecture that enables storage engines to be loaded into and unloaded from a running MySQL server.

To determine which storage engines your server supports, use the `SHOW ENGINES` statement. The value in the `Support` column indicates whether an engine can be used. A value of `YES`, `NO`, or `DEFAULT` indicates that an engine is available, not available, or available and currently set as the default storage engine.

```
mysql> SHOW ENGINES\G
*************************** 1. row ***************************
      Engine: PERFORMANCE_SCHEMA
     Support: YES
     Comment: Performance Schema
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 2. row ***************************
      Engine: InnoDB
     Support: DEFAULT
     Comment: Supports transactions, row-level locking, and foreign keys
Transactions: YES
          XA: YES
```

```
  Savepoints: YES
*************************** 3. row ***************************
      Engine: MRG_MYISAM
     Support: YES
     Comment: Collection of identical MyISAM tables
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 4. row ***************************
      Engine: BLACKHOLE
     Support: YES
     Comment: /dev/null storage engine (anything you write to it disappears)
Transactions: NO
          XA: NO
  Savepoints: NO
*************************** 5. row ***************************
      Engine: MyISAM
     Support: YES
     Comment: MyISAM storage engine
Transactions: NO
          XA: NO
  Savepoints: NO
...
```

This chapter covers use cases for special-purpose MySQL storage engines. It does not cover the default `InnoDB` storage engine or the `NDB` storage engine which are covered in Chapter 15, *The InnoDB Storage Engine* and MySQL NDB Cluster 7.5 and NDB Cluster 7.6. For advanced users, it also contains a description of the pluggable storage engine architecture (see Section 16.11, “Overview of MySQL Storage Engine Architecture”).

For information about storage engine support offered in commercial MySQL Server binaries, see *MySQL Enterprise Server 5.7*, on the MySQL Web site. The storage engines available might depend on which edition of Enterprise Server you are using.

For answers to commonly asked questions about MySQL storage engines, see Section A.2, “MySQL 8.0 FAQ: Storage Engines”.

MySQL 8.0 Supported Storage Engines

- `InnoDB`: The default storage engine in MySQL 8.0. `InnoDB` is a transaction-safe (ACID compliant) storage engine for MySQL that has commit, rollback, and crash-recovery capabilities to protect user data. `InnoDB` row-level locking (without escalation to coarser granularity locks) and Oracle-style consistent nonlocking reads increase multi-user concurrency and performance. `InnoDB` stores user data in clustered indexes to reduce I/O for common queries based on primary keys. To maintain data integrity, `InnoDB` also supports `FOREIGN KEY` referential-integrity constraints. For more information about `InnoDB`, see Chapter 15, *The InnoDB Storage Engine*.

- `MyISAM`: These tables have a small footprint. Table-level locking limits the performance in read/write workloads, so it is often used in read-only or read-mostly workloads in Web and data warehousing configurations.

- `Memory`: Stores all data in RAM, for fast access in environments that require quick lookups of non-critical data. This engine was formerly known as the `HEAP` engine. Its use cases are decreasing; `InnoDB` with its buffer pool memory area provides a general-purpose and durable way to keep most or all data in memory, and `NDBCLUSTER` provides fast key-value lookups for huge distributed data sets.

- `CSV`: Its tables are really text files with comma-separated values. CSV tables let you import or dump data in CSV format, to exchange data with scripts and applications that read and write that same format. Because CSV tables are not indexed, you typically keep the data in `InnoDB` tables during normal operation, and only use CSV tables during the import or export stage.

- `Archive`: These compact, unindexed tables are intended for storing and retrieving large amounts of seldom-referenced historical, archived, or security audit information.

- `Blackhole`: The Blackhole storage engine accepts but does not store data, similar to the Unix `/dev/null` device. Queries always return an empty set. These tables can be used in replication configurations where DML statements are sent to slave servers, but the master server does not keep its own copy of the data.
- `Merge`: Enables a MySQL DBA or developer to logically group a series of identical `MyISAM` tables and reference them as one object. Good for VLDB environments such as data warehousing.
- `Federated`: Offers the ability to link separate MySQL servers to create one logical database from many physical servers. Very good for distributed or data mart environments.
- `Example`: This engine serves as an example in the MySQL source code that illustrates how to begin writing new storage engines. It is primarily of interest to developers. The storage engine is a “stub” that does nothing. You can create tables with this engine, but no data can be stored in them or retrieved from them.

You are not restricted to using the same storage engine for an entire server or schema. You can specify the storage engine for any table. For example, an application might use mostly `InnoDB` tables, with one `CSV` table for exporting data to a spreadsheet and a few `MEMORY` tables for temporary workspaces.

Choosing a Storage Engine

The various storage engines provided with MySQL are designed with different use cases in mind. The following table provides an overview of some storage engines provided with MySQL:

Table 16.1 Storage Engines Feature Summary

Feature	MyISAM	Memory	InnoDB	Archive	NDB
Storage limits	256TB	RAM	64TB	None	384EB
Transactions	No	No	Yes	No	Yes
Locking granularity	Table	Table	Row	Row	Row
MVCC	No	No	Yes	No	No
Geospatial data type support	Yes	No	Yes	Yes	Yes
Geospatial indexing support	Yes	No	Yes[a]	No	No
B-tree indexes	Yes	Yes	Yes	No	No
T-tree indexes	No	No	No	No	Yes
Hash indexes	No	Yes	No[b]	No	Yes
Full-text search indexes	Yes	No	Yes[c]	No	No
Clustered indexes	No	No	Yes	No	No
Data caches	No	N/A	Yes	No	Yes

Feature	MyISAM	Memory	InnoDB	Archive	NDB
Index caches	Yes	N/A	Yes	No	Yes
Compressed data	Yes[d]	No	Yes[e]	Yes	No
Encrypted data[f]	Yes	Yes	Yes	Yes	Yes
Cluster database support	No	No	No	No	Yes
Replication support[g]	Yes	Yes	Yes	Yes	Yes
Foreign key support	No	No	Yes	No	Yes[h]
Backup / point-in-time recovery[i]	Yes	Yes	Yes	Yes	Yes
Query cache support	Yes	Yes	Yes	Yes	Yes
Update statistics for data dictionary	Yes	Yes	Yes	Yes	Yes

[a] InnoDB support for geospatial indexing is available in MySQL 5.7.5 and later.

[b] InnoDB utilizes hash indexes internally for its Adaptive Hash Index feature.

[c] InnoDB support for FULLTEXT indexes is available in MySQL 5.6.4 and later.

[d] Compressed MyISAM tables are supported only when using the compressed row format. Tables using the compressed row format with MyISAM are read only.

[e] Compressed InnoDB tables require the InnoDB Barracuda file format.

[f] Implemented in the server (via encryption functions). Data-at-rest tablespace encryption is available in MySQL 5.7 and later.

[g] Implemented in the server, rather than in the storage engine.

[h] Support for foreign keys is available in MySQL Cluster NDB 7.3 and later.

[i] Implemented in the server, rather than in the storage engine.

16.1 Setting the Storage Engine

When you create a new table, you can specify which storage engine to use by adding an `ENGINE` table option to the `CREATE TABLE` statement:

```
-- ENGINE=INNODB not needed unless you have set a different
-- default storage engine.
CREATE TABLE t1 (i INT) ENGINE = INNODB;
-- Simple table definitions can be switched from one to another.
CREATE TABLE t2 (i INT) ENGINE = CSV;
CREATE TABLE t3 (i INT) ENGINE = MEMORY;
```

When you omit the `ENGINE` option, the default storage engine is used. The default engine is `InnoDB` in MySQL 8.0. You can specify the default engine by using the `--default-storage-engine` server startup option, or by setting the `default-storage-engine` option in the `my.cnf` configuration file.

You can set the default storage engine for the current session by setting the `default_storage_engine` variable:

```
SET default_storage_engine=NDBCLUSTER;
```

The storage engine for `TEMPORARY` tables created with `CREATE TEMPORARY TABLE` can be set separately from the engine for permanent tables by setting the `default_tmp_storage_engine`, either at startup or at runtime.

To convert a table from one storage engine to another, use an `ALTER TABLE` statement that indicates the new engine:

```
ALTER TABLE t ENGINE = InnoDB;
```

See Section 13.1.15, “CREATE TABLE Syntax”, and Section 13.1.7, “ALTER TABLE Syntax”.

If you try to use a storage engine that is not compiled in or that is compiled in but deactivated, MySQL instead creates a table using the default storage engine. For example, in a replication setup, perhaps your master server uses `InnoDB` tables for maximum safety, but the slave servers use other storage engines for speed at the expense of durability or concurrency.

By default, a warning is generated whenever `CREATE TABLE` or `ALTER TABLE` cannot use the default storage engine. To prevent confusing, unintended behavior if the desired engine is unavailable, enable the `NO_ENGINE_SUBSTITUTION` SQL mode. If the desired engine is unavailable, this setting produces an error instead of a warning, and the table is not created or altered. See Section 5.1.8, “Server SQL Modes”.

MySQL may store a table's index and data in one or more other files, depending on the storage engine. Table and column definitions are stored in the MySQL data dictionary. Individual storage engines create any additional files required for the tables that they manage. If a table name contains special characters, the names for the table files contain encoded versions of those characters as described in Section 9.2.3, “Mapping of Identifiers to File Names”.

16.2 The MyISAM Storage Engine

`MyISAM` is based on the older (and no longer available) `ISAM` storage engine but has many useful extensions.

Table 16.2 `MyISAM` Storage Engine Features

Storage limits	256TB	**Transactions**	No	**Locking granularity**	Table
MVCC	No	**Geospatial data type support**	Yes	**Geospatial indexing support**	Yes
B-tree indexes	Yes	**T-tree indexes**	No	**Hash indexes**	No
Full-text search indexes	Yes	**Clustered indexes**	No	**Data caches**	No
Index caches	Yes	**Compressed data**	Yes[a]	**Encrypted data**[b]	Yes
Cluster database support	No	**Replication support**[c]	Yes	**Foreign key support**	No
Backup / point-in-time recovery[d]	Yes	**Query cache support**	Yes	**Update statistics for data dictionary**	Yes

[a] Compressed MyISAM tables are supported only when using the compressed row format. Tables using the compressed row format with MyISAM are read only.

[b] Implemented in the server (via encryption functions). Data-at-rest tablespace encryption is available in MySQL 5.7 and later.

[c] Implemented in the server, rather than in the storage engine.

[d] Implemented in the server, rather than in the storage engine.

Each `MyISAM` table is stored on disk in two files. The files have names that begin with the table name and have an extension to indicate the file type. The data file has an `.MYD` (`MYData`) extension.

The index file has an `.MYI` (`MYIndex`) extension. The table definition is stored in the MySQL data dictionary.

To specify explicitly that you want a `MyISAM` table, indicate that with an `ENGINE` table option:

```
CREATE TABLE t (i INT) ENGINE = MYISAM;
```

In MySQL 8.0, it is normally necessary to use `ENGINE` to specify the `MyISAM` storage engine because `InnoDB` is the default engine.

You can check or repair `MyISAM` tables with the `mysqlcheck` client or `myisamchk` utility. You can also compress `MyISAM` tables with `myisampack` to take up much less space. See Section 4.5.3, “`mysqlcheck` — A Table Maintenance Program”, Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”, and Section 4.6.6, “`myisampack` — Generate Compressed, Read-Only MyISAM Tables”.

In MySQL 8.0, the `MyISAM` storage engine provides no partitioning support. *Partitioned `MyISAM` tables created in previous versions of MySQL cannot be used in MySQL 8.0*. For more information, see Section 22.6.2, “Partitioning Limitations Relating to Storage Engines”. For help with upgrading such tables so that they can be used in MySQL 8.0, see Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0”.

`MyISAM` tables have the following characteristics:

- All data values are stored with the low byte first. This makes the data machine and operating system independent. The only requirements for binary portability are that the machine uses two's-complement signed integers and IEEE floating-point format. These requirements are widely used among mainstream machines. Binary compatibility might not be applicable to embedded systems, which sometimes have peculiar processors.

  There is no significant speed penalty for storing data low byte first; the bytes in a table row normally are unaligned and it takes little more processing to read an unaligned byte in order than in reverse order. Also, the code in the server that fetches column values is not time critical compared to other code.

- All numeric key values are stored with the high byte first to permit better index compression.

- Large files (up to 63-bit file length) are supported on file systems and operating systems that support large files.

- There is a limit of $(2^{32})^2$ (1.844E+19) rows in a `MyISAM` table.

- The maximum number of indexes per `MyISAM` table is 64.

  The maximum number of columns per index is 16.

- The maximum key length is 1000 bytes. This can also be changed by changing the source and recompiling. For the case of a key longer than 250 bytes, a larger key block size than the default of 1024 bytes is used.

- When rows are inserted in sorted order (as when you are using an `AUTO_INCREMENT` column), the index tree is split so that the high node only contains one key. This improves space utilization in the index tree.

- Internal handling of one `AUTO_INCREMENT` column per table is supported. `MyISAM` automatically updates this column for `INSERT` and `UPDATE` operations. This makes `AUTO_INCREMENT` columns faster (at least 10%). Values at the top of the sequence are not reused after being deleted. (When an `AUTO_INCREMENT` column is defined as the last column of a multiple-column index, reuse of values deleted from the top of a sequence does occur.) The `AUTO_INCREMENT` value can be reset with `ALTER TABLE` or `myisamchk`.

- Dynamic-sized rows are much less fragmented when mixing deletes with updates and inserts. This is done by automatically combining adjacent deleted blocks and by extending blocks if the next block is deleted.
- `MyISAM` supports concurrent inserts: If a table has no free blocks in the middle of the data file, you can `INSERT` new rows into it at the same time that other threads are reading from the table. A free block can occur as a result of deleting rows or an update of a dynamic length row with more data than its current contents. When all free blocks are used up (filled in), future inserts become concurrent again. See Section 8.11.3, “Concurrent Inserts”.
- You can put the data file and index file in different directories on different physical devices to get more speed with the `DATA DIRECTORY` and `INDEX DIRECTORY` table options to `CREATE TABLE`. See Section 13.1.15, “CREATE TABLE Syntax”.
- `BLOB` and `TEXT` columns can be indexed.
- `NULL` values are permitted in indexed columns. This takes 0 to 1 bytes per key.
- Each character column can have a different character set. See Section 10.1, “Character Set Support”.
- There is a flag in the `MyISAM` index file that indicates whether the table was closed correctly. If `mysqld` is started with the `--myisam-recover-options` option, `MyISAM` tables are automatically checked when opened, and are repaired if the table wasn't closed properly.
- `myisamchk` marks tables as checked if you run it with the `--update-state` option. `myisamchk --fast` checks only those tables that don't have this mark.
- `myisamchk --analyze` stores statistics for portions of keys, as well as for entire keys.
- `myisampack` can pack `BLOB` and `VARCHAR` columns.

`MyISAM` also supports the following features:

- Support for a true `VARCHAR` type; a `VARCHAR` column starts with a length stored in one or two bytes.
- Tables with `VARCHAR` columns may have fixed or dynamic row length.
- The sum of the lengths of the `VARCHAR` and `CHAR` columns in a table may be up to 64KB.
- Arbitrary length `UNIQUE` constraints.

Additional Resources

- A forum dedicated to the `MyISAM` storage engine is available at http://forums.mysql.com/list.php?21.

16.2.1 MyISAM Startup Options

The following options to `mysqld` can be used to change the behavior of `MyISAM` tables. For additional information, see Section 5.1.4, “Server Command Options”.

Table 16.3 `MyISAM` Option/Variable Reference

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
bulk_insert_buffer_size	Yes	Yes	Yes		Both	Yes
concurrent_insert	Yes	Yes	Yes		Global	Yes
delay-key-write	Yes	Yes			Global	Yes
- *Variable*: delay_key_write			Yes		Global	Yes
have_rtree_keys			Yes		Global	No
key_buffer_size	Yes	Yes	Yes		Global	Yes

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
log-isam	Yes	Yes				
myisam-block-size	Yes	Yes				
myisam_data_pointer_size	Yes	Yes	Yes		Global	Yes
myisam_max_sort_file_size	Yes	Yes	Yes		Global	Yes
myisam_mmap_size	Yes	Yes	Yes		Global	No
myisam-recover-options	Yes	Yes				
- *Variable*: myisam_recover_options						
myisam_recover_options			Yes		Global	No
myisam_repair_threads	Yes	Yes	Yes		Both	Yes
myisam_sort_buffer_size	Yes	Yes	Yes		Both	Yes
myisam_stats_method	Yes	Yes	Yes		Both	Yes
myisam_use_mmap	Yes	Yes	Yes		Global	Yes
skip-concurrent-insert	Yes	Yes				
- *Variable*: concurrent_insert						
tmp_table_size	Yes	Yes	Yes		Both	Yes

- `--myisam-recover-options=`*mode*

  Set the mode for automatic recovery of crashed `MyISAM` tables.

- `--delay-key-write=ALL`

  Don't flush key buffers between writes for any `MyISAM` table.

Note

If you do this, you should not access `MyISAM` tables from another program (such as from another MySQL server or with `myisamchk`) when the tables are in use. Doing so risks index corruption. Using `--external-locking` does not eliminate this risk.

The following system variables affect the behavior of `MyISAM` tables. For additional information, see Section 5.1.5, “Server System Variables”.

- `bulk_insert_buffer_size`

  The size of the tree cache used in bulk insert optimization.

Note

This is a limit *per thread*!

- `myisam_max_sort_file_size`

  The maximum size of the temporary file that MySQL is permitted to use while re-creating a `MyISAM` index (during `REPAIR TABLE`, `ALTER TABLE`, or `LOAD DATA INFILE`). If the file size would be larger than this value, the index is created using the key cache instead, which is slower. The value is given in bytes.

- `myisam_sort_buffer_size`

  Set the size of the buffer used when recovering tables.

Automatic recovery is activated if you start `mysqld` with the `--myisam-recover-options` option. In this case, when the server opens a `MyISAM` table, it checks whether the table is marked as crashed or whether the open count variable for the table is not 0 and you are running the server with external locking disabled. If either of these conditions is true, the following happens:

- The server checks the table for errors.
- If the server finds an error, it tries to do a fast table repair (with sorting and without re-creating the data file).
- If the repair fails because of an error in the data file (for example, a duplicate-key error), the server tries again, this time re-creating the data file.
- If the repair still fails, the server tries once more with the old repair option method (write row by row without sorting). This method should be able to repair any type of error and has low disk space requirements.

If the recovery wouldn't be able to recover all rows from previously completed statements and you didn't specify `FORCE` in the value of the `--myisam-recover-options` option, automatic repair aborts with an error message in the error log:

```
Error: Couldn't repair table: test.g00pages
```

If you specify `FORCE`, a warning like this is written instead:

```
Warning: Found 344 of 354 rows when repairing ./test/g00pages
```

If the automatic recovery value includes `BACKUP`, the recovery process creates files with names of the form *`tbl_name-datetime`*`.BAK`. You should have a `cron` script that automatically moves these files from the database directories to backup media.

16.2.2 Space Needed for Keys

`MyISAM` tables use B-tree indexes. You can roughly calculate the size for the index file as `(key_length+4)/0.67`, summed over all keys. This is for the worst case when all keys are inserted in sorted order and the table doesn't have any compressed keys.

String indexes are space compressed. If the first index part is a string, it is also prefix compressed. Space compression makes the index file smaller than the worst-case figure if a string column has a lot of trailing space or is a `VARCHAR` column that is not always used to the full length. Prefix compression is used on keys that start with a string. Prefix compression helps if there are many strings with an identical prefix.

In `MyISAM` tables, you can also prefix compress numbers by specifying the `PACK_KEYS=1` table option when you create the table. Numbers are stored with the high byte first, so this helps when you have many integer keys that have an identical prefix.

16.2.3 MyISAM Table Storage Formats

`MyISAM` supports three different storage formats. Two of them, fixed and dynamic format, are chosen automatically depending on the type of columns you are using. The third, compressed format, can be created only with the `myisampack` utility (see Section 4.6.6, “`myisampack` — Generate Compressed, Read-Only MyISAM Tables”).

When you use `CREATE TABLE` or `ALTER TABLE` for a table that has no `BLOB` or `TEXT` columns, you can force the table format to `FIXED` or `DYNAMIC` with the `ROW_FORMAT` table option.

See Section 13.1.15, “CREATE TABLE Syntax”, for information about ROW_FORMAT.

You can decompress (unpack) compressed MyISAM tables using myisamchk --unpack; see Section 4.6.4, “myisamchk — MyISAM Table-Maintenance Utility”, for more information.

16.2.3.1 Static (Fixed-Length) Table Characteristics

Static format is the default for MyISAM tables. It is used when the table contains no variable-length columns (VARCHAR, VARBINARY, BLOB, or TEXT). Each row is stored using a fixed number of bytes.

Of the three MyISAM storage formats, static format is the simplest and most secure (least subject to corruption). It is also the fastest of the on-disk formats due to the ease with which rows in the data file can be found on disk: To look up a row based on a row number in the index, multiply the row number by the row length to calculate the row position. Also, when scanning a table, it is very easy to read a constant number of rows with each disk read operation.

The security is evidenced if your computer crashes while the MySQL server is writing to a fixed-format MyISAM file. In this case, myisamchk can easily determine where each row starts and ends, so it can usually reclaim all rows except the partially written one. MyISAM table indexes can always be reconstructed based on the data rows.

Note

Fixed-length row format is only available for tables without BLOB or TEXT columns. Creating a table with these columns with an explicit ROW_FORMAT clause will not raise an error or warning; the format specification will be ignored.

Static-format tables have these characteristics:

- CHAR and VARCHAR columns are space-padded to the specified column width, although the column type is not altered. BINARY and VARBINARY columns are padded with 0x00 bytes to the column width.
- NULL columns require additional space in the row to record whether their values are NULL. Each NULL column takes one bit extra, rounded up to the nearest byte.
- Very quick.
- Easy to cache.
- Easy to reconstruct after a crash, because rows are located in fixed positions.
- Reorganization is unnecessary unless you delete a huge number of rows and want to return free disk space to the operating system. To do this, use OPTIMIZE TABLE or myisamchk -r.
- Usually require more disk space than dynamic-format tables.
- The expected row length in bytes for static-sized rows is calculated using the following expression:

```
row length = 1
             + (sum of column lengths)
             + (number of NULL columns + delete_flag + 7)/8
             + (number of variable-length columns)
```

delete_flag is 1 for tables with static row format. Static tables use a bit in the row record for a flag that indicates whether the row has been deleted. *delete_flag* is 0 for dynamic tables because the flag is stored in the dynamic row header.

16.2.3.2 Dynamic Table Characteristics

Dynamic storage format is used if a MyISAM table contains any variable-length columns (VARCHAR, VARBINARY, BLOB, or TEXT), or if the table was created with the ROW_FORMAT=DYNAMIC table option.

Dynamic format is a little more complex than static format because each row has a header that indicates how long it is. A row can become fragmented (stored in noncontiguous pieces) when it is made longer as a result of an update.

You can use `OPTIMIZE TABLE` or `myisamchk -r` to defragment a table. If you have fixed-length columns that you access or change frequently in a table that also contains some variable-length columns, it might be a good idea to move the variable-length columns to other tables just to avoid fragmentation.

Dynamic-format tables have these characteristics:

- All string columns are dynamic except those with a length less than four.
- Each row is preceded by a bitmap that indicates which columns contain the empty string (for string columns) or zero (for numeric columns). This does not include columns that contain `NULL` values. If a string column has a length of zero after trailing space removal, or a numeric column has a value of zero, it is marked in the bitmap and not saved to disk. Nonempty strings are saved as a length byte plus the string contents.
- `NULL` columns require additional space in the row to record whether their values are `NULL`. Each `NULL` column takes one bit extra, rounded up to the nearest byte.
- Much less disk space usually is required than for fixed-length tables.
- Each row uses only as much space as is required. However, if a row becomes larger, it is split into as many pieces as are required, resulting in row fragmentation. For example, if you update a row with information that extends the row length, the row becomes fragmented. In this case, you may have to run `OPTIMIZE TABLE` or `myisamchk -r` from time to time to improve performance. Use `myisamchk -ei` to obtain table statistics.
- More difficult than static-format tables to reconstruct after a crash, because rows may be fragmented into many pieces and links (fragments) may be missing.
- The expected row length for dynamic-sized rows is calculated using the following expression:

```
3
+ (number of columns + 7) / 8
+ (number of char columns)
+ (packed size of numeric columns)
+ (length of strings)
+ (number of NULL columns + 7) / 8
```

  There is a penalty of 6 bytes for each link. A dynamic row is linked whenever an update causes an enlargement of the row. Each new link is at least 20 bytes, so the next enlargement probably goes in the same link. If not, another link is created. You can find the number of links using `myisamchk -ed`. All links may be removed with `OPTIMIZE TABLE` or `myisamchk -r`.

16.2.3.3 Compressed Table Characteristics

Compressed storage format is a read-only format that is generated with the `myisampack` tool. Compressed tables can be uncompressed with `myisamchk`.

Compressed tables have the following characteristics:

- Compressed tables take very little disk space. This minimizes disk usage, which is helpful when using slow disks (such as CD-ROMs).
- Each row is compressed separately, so there is very little access overhead. The header for a row takes up one to three bytes depending on the biggest row in the table. Each column is compressed differently. There is usually a different Huffman tree for each column. Some of the compression types are:

  - Suffix space compression.
  - Prefix space compression.
  - Numbers with a value of zero are stored using one bit.
  - If values in an integer column have a small range, the column is stored using the smallest possible type. For example, a `BIGINT` column (eight bytes) can be stored as a `TINYINT` column (one byte) if all its values are in the range from `-128` to `127`.
  - If a column has only a small set of possible values, the data type is converted to `ENUM`.
  - A column may use any combination of the preceding compression types.
- Can be used for fixed-length or dynamic-length rows.

Note

While a compressed table is read only, and you cannot therefore update or add rows in the table, DDL (Data Definition Language) operations are still valid. For example, you may still use `DROP` to drop the table, and `TRUNCATE TABLE` to empty the table.

16.2.4 MyISAM Table Problems

The file format that MySQL uses to store data has been extensively tested, but there are always circumstances that may cause database tables to become corrupted. The following discussion describes how this can happen and how to handle it.

16.2.4.1 Corrupted MyISAM Tables

Even though the `MyISAM` table format is very reliable (all changes to a table made by an SQL statement are written before the statement returns), you can still get corrupted tables if any of the following events occur:

- The `mysqld` process is killed in the middle of a write.
- An unexpected computer shutdown occurs (for example, the computer is turned off).
- Hardware failures.
- You are using an external program (such as `myisamchk`) to modify a table that is being modified by the server at the same time.
- A software bug in the MySQL or `MyISAM` code.

Typical symptoms of a corrupt table are:

- You get the following error while selecting data from the table:

```
Incorrect key file for table: '...'. Try to repair it
```

- Queries don't find rows in the table or return incomplete results.

You can check the health of a `MyISAM` table using the `CHECK TABLE` statement, and repair a corrupted `MyISAM` table with `REPAIR TABLE`. When `mysqld` is not running, you can also check or repair a table with the `myisamchk` command. See Section 13.7.2.2, “CHECK TABLE Syntax”, Section 13.7.2.5, “REPAIR TABLE Syntax”, and Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”.

If your tables become corrupted frequently, you should try to determine why this is happening. The most important thing to know is whether the table became corrupted as a result of a server crash. You can verify this easily by looking for a recent `restarted mysqld` message in the error log. If there is such a message, it is likely that table corruption is a result of the server dying. Otherwise, corruption may have occurred during normal operation. This is a bug. You should try to create a reproducible test case that demonstrates the problem. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”, and Section 28.5, “Debugging and Porting MySQL”.

16.2.4.2 Problems from Tables Not Being Closed Properly

Each `MyISAM` index file (`.MYI` file) has a counter in the header that can be used to check whether a table has been closed properly. If you get the following warning from `CHECK TABLE` or `myisamchk`, it means that this counter has gone out of sync:

```
clients are using or haven't closed the table properly
```

This warning doesn't necessarily mean that the table is corrupted, but you should at least check the table.

The counter works as follows:

- The first time a table is updated in MySQL, a counter in the header of the index files is incremented.
- The counter is not changed during further updates.
- When the last instance of a table is closed (because a `FLUSH TABLES` operation was performed or because there is no room in the table cache), the counter is decremented if the table has been updated at any point.
- When you repair the table or check the table and it is found to be okay, the counter is reset to zero.
- To avoid problems with interaction with other processes that might check the table, the counter is not decremented on close if it was zero.

In other words, the counter can become incorrect only under these conditions:

- A `MyISAM` table is copied without first issuing `LOCK TABLES` and `FLUSH TABLES`.
- MySQL has crashed between an update and the final close. (The table may still be okay because MySQL always issues writes for everything between each statement.)
- A table was modified by `myisamchk --recover` or `myisamchk --update-state` at the same time that it was in use by `mysqld`.
- Multiple `mysqld` servers are using the table and one server performed a `REPAIR TABLE` or `CHECK TABLE` on the table while it was in use by another server. In this setup, it is safe to use `CHECK TABLE`, although you might get the warning from other servers. However, `REPAIR TABLE` should be avoided because when one server replaces the data file with a new one, this is not known to the other servers.

  In general, it is a bad idea to share a data directory among multiple servers. See Section 5.7, “Running Multiple MySQL Instances on One Machine”, for additional discussion.

16.3 The MEMORY Storage Engine

The `MEMORY` storage engine (formerly known as `HEAP`) creates special-purpose tables with contents that are stored in memory. Because the data is vulnerable to crashes, hardware issues, or power outages, only use these tables as temporary work areas or read-only caches for data pulled from other tables.

Table 16.4 `MEMORY` Storage Engine Features

Storage limits	RAM	**Transactions**	No	**Locking granularity**	Table
MVCC	No	**Geospatial data type support**	No	**Geospatial indexing support**	No
B-tree indexes	Yes	**T-tree indexes**	No	**Hash indexes**	Yes
Full-text search indexes	No	**Clustered indexes**	No	**Data caches**	N/A
Index caches	N/A	**Compressed data**	No	**Encrypted data**[a]	Yes
Cluster database support	No	**Replication support**[b]	Yes	**Foreign key support**	No
Backup / point-in-time recovery[c]	Yes	**Query cache support**	Yes	**Update statistics for data dictionary**	Yes

[a] Implemented in the server (via encryption functions). Data-at-rest tablespace encryption is available in MySQL 5.7 and later.
[b] Implemented in the server, rather than in the storage engine.
[c] Implemented in the server, rather than in the storage engine.

When to Use MEMORY or MySQL Cluster. Developers looking to deploy applications that use the `MEMORY` storage engine for important, highly available, or frequently updated data should consider whether MySQL Cluster is a better choice. A typical use case for the `MEMORY` engine involves these characteristics:

- Operations involving transient, non-critical data such as session management or caching. When the MySQL server halts or restarts, the data in `MEMORY` tables is lost.
- In-memory storage for fast access and low latency. Data volume can fit entirely in memory without causing the operating system to swap out virtual memory pages.
- A read-only or read-mostly data access pattern (limited updates).

MySQL Cluster offers the same features as the `MEMORY` engine with higher performance levels, and provides additional features not available with `MEMORY`:

- Row-level locking and multiple-thread operation for low contention between clients.
- Scalability even with statement mixes that include writes.
- Optional disk-backed operation for data durability.
- Shared-nothing architecture and multiple-host operation with no single point of failure, enabling 99.999% availability.
- Automatic data distribution across nodes; application developers need not craft custom sharding or partitioning solutions.
- Support for variable-length data types (including `BLOB` and `TEXT`) not supported by `MEMORY`.

`MEMORY` tables cannot be partitioned.

Performance Characteristics

`MEMORY` performance is constrained by contention resulting from single-thread execution and table lock overhead when processing updates. This limits scalability when load increases, particularly for statement mixes that include writes.

Despite the in-memory processing for `MEMORY` tables, they are not necessarily faster than `InnoDB` tables on a busy server, for general-purpose queries, or under a read/write workload. In particular, the table locking involved with performing updates can slow down concurrent usage of `MEMORY` tables from multiple sessions.

Depending on the kinds of queries performed on a MEMORY table, you might create indexes as either the default hash data structure (for looking up single values based on a unique key), or a general-purpose B-tree data structure (for all kinds of queries involving equality, inequality, or range operators such as less than or greater than). The following sections illustrate the syntax for creating both kinds of indexes. A common performance issue is using the default hash indexes in workloads where B-tree indexes are more efficient.

Characteristics of MEMORY Tables

The MEMORY storage engine does not create any files on disk. The table definition is stored in the MySQL data dictionary.

MEMORY tables have the following characteristics:

- Space for MEMORY tables is allocated in small blocks. Tables use 100% dynamic hashing for inserts. No overflow area or extra key space is needed. No extra space is needed for free lists. Deleted rows are put in a linked list and are reused when you insert new data into the table. MEMORY tables also have none of the problems commonly associated with deletes plus inserts in hashed tables.
- MEMORY tables use a fixed-length row-storage format. Variable-length types such as VARCHAR are stored using a fixed length.
- MEMORY tables cannot contain BLOB or TEXT columns.
- MEMORY includes support for AUTO_INCREMENT columns.
- Non-TEMPORARY MEMORY tables are shared among all clients, just like any other non-TEMPORARY table.

DDL Operations for MEMORY Tables

To create a MEMORY table, specify the clause ENGINE=MEMORY on the CREATE TABLE statement.

```
CREATE TABLE t (i INT) ENGINE = MEMORY;
```

As indicated by the engine name, MEMORY tables are stored in memory. They use hash indexes by default, which makes them very fast for single-value lookups, and very useful for creating temporary tables. However, when the server shuts down, all rows stored in MEMORY tables are lost. The tables themselves continue to exist because their definitions are stored in MySQL data dictionary, but they are empty when the server restarts.

This example shows how you might create, use, and remove a MEMORY table:

```
mysql> CREATE TABLE test ENGINE=MEMORY
    ->     SELECT ip,SUM(downloads) AS down
    ->     FROM log_table GROUP BY ip;
mysql> SELECT COUNT(ip),AVG(down) FROM test;
mysql> DROP TABLE test;
```

The maximum size of MEMORY tables is limited by the max_heap_table_size system variable, which has a default value of 16MB. To enforce different size limits for MEMORY tables, change the value of this variable. The value in effect for CREATE TABLE, or a subsequent ALTER TABLE or TRUNCATE TABLE, is the value used for the life of the table. A server restart also sets the maximum size of existing MEMORY tables to the global max_heap_table_size value. You can set the size for individual tables as described later in this section.

Indexes

The MEMORY storage engine supports both HASH and BTREE indexes. You can specify one or the other for a given index by adding a USING clause as shown here:

```
CREATE TABLE lookup
    (id INT, INDEX USING HASH (id))
    ENGINE = MEMORY;
CREATE TABLE lookup
    (id INT, INDEX USING BTREE (id))
    ENGINE = MEMORY;
```

For general characteristics of B-tree and hash indexes, see Section 8.3.1, “How MySQL Uses Indexes”.

`MEMORY` tables can have up to 64 indexes per table, 16 columns per index and a maximum key length of 3072 bytes.

If a `MEMORY` table hash index has a high degree of key duplication (many index entries containing the same value), updates to the table that affect key values and all deletes are significantly slower. The degree of this slowdown is proportional to the degree of duplication (or, inversely proportional to the index cardinality). You can use a `BTREE` index to avoid this problem.

`MEMORY` tables can have nonunique keys. (This is an uncommon feature for implementations of hash indexes.)

Columns that are indexed can contain `NULL` values.

User-Created and Temporary Tables

`MEMORY` table contents are stored in memory, which is a property that `MEMORY` tables share with internal temporary tables that the server creates on the fly while processing queries. However, the two types of tables differ in that `MEMORY` tables are not subject to storage conversion, whereas internal temporary tables are:

- If an internal temporary table becomes too large, the server automatically converts it to on-disk storage, as described in Section 8.4.4, “Internal Temporary Table Use in MySQL”.
- User-created `MEMORY` tables are never converted to disk tables.

Loading Data

To populate a `MEMORY` table when the MySQL server starts, you can use the `--init-file` option. For example, you can put statements such as `INSERT INTO ... SELECT` or `LOAD DATA INFILE` into this file to load the table from a persistent data source. See Section 5.1.4, “Server Command Options”, and Section 13.2.7, “LOAD DATA INFILE Syntax”.

MEMORY Tables and Replication

A server's `MEMORY` tables become empty when it is shut down and restarted. If the server is a replication master, its slaves are not aware that these tables have become empty, so you see out-of-date content if you select data from the tables on the slaves. To synchronize master and slave `MEMORY` tables, when a `MEMORY` table is used on a master for the first time since it was started, a `DELETE` statement is written to the master's binary log, to empty the table on the slaves also. The slave still has outdated data in the table during the interval between the master's restart and its first use of the table. To avoid this interval when a direct query to the slave could return stale data, use the `--init-file` option to populate the `MEMORY` table on the master at startup.

Managing Memory Use

The server needs sufficient memory to maintain all `MEMORY` tables that are in use at the same time.

Memory is not reclaimed if you delete individual rows from a `MEMORY` table. Memory is reclaimed only when the entire table is deleted. Memory that was previously used for deleted rows is re-used for new rows within the same table. To free all the memory used by a `MEMORY` table when you no longer require its contents, execute `DELETE` or `TRUNCATE TABLE` to remove all rows, or remove the table

altogether using `DROP TABLE`. To free up the memory used by deleted rows, use `ALTER TABLE ENGINE=MEMORY` to force a table rebuild.

The memory needed for one row in a `MEMORY` table is calculated using the following expression:

```
SUM_OVER_ALL_BTREE_KEYS(max_length_of_key + sizeof(char*) * 4)
+ SUM_OVER_ALL_HASH_KEYS(sizeof(char*) * 2)
+ ALIGN(length_of_row+1, sizeof(char*))
```

`ALIGN()` represents a round-up factor to cause the row length to be an exact multiple of the `char` pointer size. `sizeof(char*)` is 4 on 32-bit machines and 8 on 64-bit machines.

As mentioned earlier, the `max_heap_table_size` system variable sets the limit on the maximum size of `MEMORY` tables. To control the maximum size for individual tables, set the session value of this variable before creating each table. (Do not change the global `max_heap_table_size` value unless you intend the value to be used for `MEMORY` tables created by all clients.) The following example creates two `MEMORY` tables, with a maximum size of 1MB and 2MB, respectively:

```
mysql> SET max_heap_table_size = 1024*1024;
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t1 (id INT, UNIQUE(id)) ENGINE = MEMORY;
Query OK, 0 rows affected (0.01 sec)

mysql> SET max_heap_table_size = 1024*1024*2;
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE TABLE t2 (id INT, UNIQUE(id)) ENGINE = MEMORY;
Query OK, 0 rows affected (0.00 sec)
```

Both tables revert to the server's global `max_heap_table_size` value if the server restarts.

You can also specify a `MAX_ROWS` table option in `CREATE TABLE` statements for `MEMORY` tables to provide a hint about the number of rows you plan to store in them. This does not enable the table to grow beyond the `max_heap_table_size` value, which still acts as a constraint on maximum table size. For maximum flexibility in being able to use `MAX_ROWS`, set `max_heap_table_size` at least as high as the value to which you want each `MEMORY` table to be able to grow.

Additional Resources

A forum dedicated to the `MEMORY` storage engine is available at http://forums.mysql.com/list.php?92.

16.4 The CSV Storage Engine

The `CSV` storage engine stores data in text files using comma-separated values format.

The `CSV` storage engine is always compiled into the MySQL server.

To examine the source for the `CSV` engine, look in the `storage/csv` directory of a MySQL source distribution.

When you create a `CSV` table, the server creates a data file. The data file name begins with the table name and has a `.CSV` extension. The data file is a plain text file. When you store data into the table, the storage engine saves it into the data file in comma-separated values format.

```
mysql> CREATE TABLE test (i INT NOT NULL, c CHAR(10) NOT NULL)
    -> ENGINE = CSV;
Query OK, 0 rows affected (0.12 sec)

mysql> INSERT INTO test VALUES(1,'record one'),(2,'record two');
Query OK, 2 rows affected (0.00 sec)
Records: 2  Duplicates: 0  Warnings: 0
```

```
mysql> SELECT * FROM test;
+------+------------+
| i    | c          |
+------+------------+
|    1 | record one |
|    2 | record two |
+------+------------+
2 rows in set (0.00 sec)
```

Creating a CSV table also creates a corresponding Metafile that stores the state of the table and the number of rows that exist in the table. The name of this file is the same as the name of the table with the extension `CSM`.

If you examine the `test.CSV` file in the database directory created by executing the preceding statements, its contents should look like this:

```
"1","record one"
"2","record two"
```

This format can be read, and even written, by spreadsheet applications such as Microsoft Excel or StarOffice Calc.

16.4.1 Repairing and Checking CSV Tables

The CSV storage engines supports the `CHECK` and `REPAIR` statements to verify and if possible repair a damaged CSV table.

When running the `CHECK` statement, the CSV file will be checked for validity by looking for the correct field separators, escaped fields (matching or missing quotation marks), the correct number of fields compared to the table definition and the existence of a corresponding CSV metafile. The first invalid row discovered will report an error. Checking a valid table produces output like that shown below:

```
mysql> check table csvtest;
+--------------+-------+----------+----------+
| Table        | Op    | Msg_type | Msg_text |
+--------------+-------+----------+----------+
| test.csvtest | check | status   | OK       |
+--------------+-------+----------+----------+
1 row in set (0.00 sec)
```

A check on a corrupted table returns a fault:

```
mysql> check table csvtest;
+--------------+-------+----------+----------+
| Table        | Op    | Msg_type | Msg_text |
+--------------+-------+----------+----------+
| test.csvtest | check | error    | Corrupt  |
+--------------+-------+----------+----------+
1 row in set (0.01 sec)
```

If the check fails, the table is marked as crashed (corrupt). Once a table has been marked as corrupt, it is automatically repaired when you next run `CHECK` or execute a `SELECT` statement. The corresponding corrupt status and new status will be displayed when running `CHECK`:

```
mysql> check table csvtest;
+--------------+-------+----------+---------------------------+
| Table        | Op    | Msg_type | Msg_text                  |
+--------------+-------+----------+---------------------------+
| test.csvtest | check | warning  | Table is marked as crashed |
| test.csvtest | check | status   | OK                        |
+--------------+-------+----------+---------------------------+
2 rows in set (0.08 sec)
```

To repair a table you can use `REPAIR`, this copies as many valid rows from the existing CSV data as possible, and then replaces the existing CSV file with the recovered rows. Any rows beyond the corrupted data are lost.

```
mysql> repair table csvtest;
+-------------+--------+----------+----------+
| Table       | Op     | Msg_type | Msg_text |
+-------------+--------+----------+----------+
| test.csvtest | repair | status   | OK       |
+-------------+--------+----------+----------+
1 row in set (0.02 sec)
```

Warning

During repair, only the rows from the CSV file up to the first damaged row are copied to the new table. All other rows from the first damaged row to the end of the table are removed, even valid rows.

16.4.2 CSV Limitations

The `CSV` storage engine does not support indexing.

The `CSV` storage engine does not support partitioning.

All tables that you create using the `CSV` storage engine must have the `NOT NULL` attribute on all columns. However, for backward compatibility, you can continue to use tables with nullable columns that were created in previous MySQL releases. (Bug #32050)

16.5 The ARCHIVE Storage Engine

The `ARCHIVE` storage engine produces special-purpose tables that store large amounts of unindexed data in a very small footprint.

Table 16.5 `ARCHIVE` Storage Engine Features

Storage limits	None	**Transactions**	No	**Locking granularity**	Row
MVCC	No	**Geospatial data type support**	Yes	**Geospatial indexing support**	No
B-tree indexes	No	**T-tree indexes**	No	**Hash indexes**	No
Full-text search indexes	No	**Clustered indexes**	No	**Data caches**	No
Index caches	No	**Compressed data**	Yes	**Encrypted data**[a]	Yes
Cluster database support	No	**Replication support**[b]	Yes	**Foreign key support**	No
Backup / point-in-time recovery[c]	Yes	**Query cache support**	Yes	**Update statistics for data dictionary**	Yes

[a] Implemented in the server (via encryption functions). Data-at-rest tablespace encryption is available in MySQL 5.7 and later.

[b] Implemented in the server, rather than in the storage engine.

[c] Implemented in the server, rather than in the storage engine.

The `ARCHIVE` storage engine is included in MySQL binary distributions. To enable this storage engine if you build MySQL from source, invoke `CMake` with the `-DWITH_ARCHIVE_STORAGE_ENGINE` option.

To examine the source for the `ARCHIVE` engine, look in the `storage/archive` directory of a MySQL source distribution.

You can check whether the `ARCHIVE` storage engine is available with the `SHOW ENGINES` statement.

When you create an `ARCHIVE` table, the storage engine creates files with names that begin with the table name. The data file has an extension of `.ARZ`. An `.ARN` file may appear during optimization operations.

The ARCHIVE engine supports INSERT, REPLACE, and SELECT, but not DELETE or UPDATE. It does support ORDER BY operations, BLOB columns, and basically all but spatial data types (see Section 11.5.1, “Spatial Data Types”). The ARCHIVE engine uses row-level locking.

The ARCHIVE engine supports the AUTO_INCREMENT column attribute. The AUTO_INCREMENT column can have either a unique or nonunique index. Attempting to create an index on any other column results in an error. The ARCHIVE engine also supports the AUTO_INCREMENT table option in CREATE TABLE statements to specify the initial sequence value for a new table or reset the sequence value for an existing table, respectively.

ARCHIVE does not support inserting a value into an AUTO_INCREMENT column less than the current maximum column value. Attempts to do so result in an ER_DUP_KEY error.

The ARCHIVE engine ignores BLOB columns if they are not requested and scans past them while reading.

The ARCHIVE storage engine does not support partitioning.

Storage: Rows are compressed as they are inserted. The ARCHIVE engine uses zlib lossless data compression (see http://www.zlib.net/). You can use OPTIMIZE TABLE to analyze the table and pack it into a smaller format (for a reason to use OPTIMIZE TABLE, see later in this section). The engine also supports CHECK TABLE. There are several types of insertions that are used:

- An INSERT statement just pushes rows into a compression buffer, and that buffer flushes as necessary. The insertion into the buffer is protected by a lock. A SELECT forces a flush to occur.
- A bulk insert is visible only after it completes, unless other inserts occur at the same time, in which case it can be seen partially. A SELECT never causes a flush of a bulk insert unless a normal insert occurs while it is loading.

Retrieval: On retrieval, rows are uncompressed on demand; there is no row cache. A SELECT operation performs a complete table scan: When a SELECT occurs, it finds out how many rows are currently available and reads that number of rows. SELECT is performed as a consistent read. Note that lots of SELECT statements during insertion can deteriorate the compression, unless only bulk or delayed inserts are used. To achieve better compression, you can use OPTIMIZE TABLE or REPAIR TABLE. The number of rows in ARCHIVE tables reported by SHOW TABLE STATUS is always accurate. See Section 13.7.2.4, “OPTIMIZE TABLE Syntax”, Section 13.7.2.5, “REPAIR TABLE Syntax”, and Section 13.7.5.36, “SHOW TABLE STATUS Syntax”.

Additional Resources

- A forum dedicated to the ARCHIVE storage engine is available at http://forums.mysql.com/list.php?112.

16.6 The BLACKHOLE Storage Engine

The BLACKHOLE storage engine acts as a “black hole” that accepts data but throws it away and does not store it. Retrievals always return an empty result:

```
mysql> CREATE TABLE test(i INT, c CHAR(10)) ENGINE = BLACKHOLE;
Query OK, 0 rows affected (0.03 sec)

mysql> INSERT INTO test VALUES(1,'record one'),(2,'record two');
Query OK, 2 rows affected (0.00 sec)
Records: 2  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM test;
Empty set (0.00 sec)
```

To enable the BLACKHOLE storage engine if you build MySQL from source, invoke CMake with the -DWITH_BLACKHOLE_STORAGE_ENGINE option.

To examine the source for the BLACKHOLE engine, look in the sql directory of a MySQL source distribution.

When you create a BLACKHOLE table, the server creates the table definition in the global data dictionary. There are no files associated with the table.

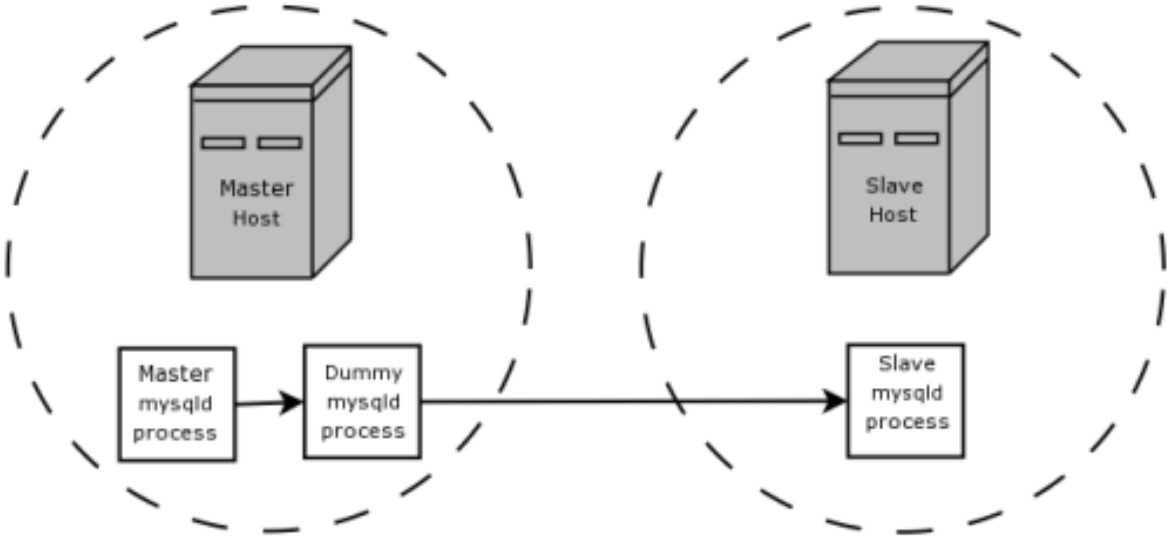
The BLACKHOLE storage engine supports all kinds of indexes. That is, you can include index declarations in the table definition.

The BLACKHOLE storage engine does not support partitioning.

You can check whether the BLACKHOLE storage engine is available with the SHOW ENGINES statement.

Inserts into a BLACKHOLE table do not store any data, but if statement based binary logging is enabled, the SQL statements are logged and replicated to slave servers. This can be useful as a repeater or filter mechanism.

Suppose that your application requires slave-side filtering rules, but transferring all binary log data to the slave first results in too much traffic. In such a case, it is possible to set up on the master host a "dummy" slave process whose default storage engine is BLACKHOLE, depicted as follows:



The master writes to its binary log. The "dummy" mysqld process acts as a slave, applying the desired combination of replicate-do-* and replicate-ignore-* rules, and writes a new, filtered binary log of its own. (See Section 18.1.6, "Replication and Binary Logging Options and Variables".) This filtered log is provided to the slave.

The dummy process does not actually store any data, so there is little processing overhead incurred by running the additional mysqld process on the replication master host. This type of setup can be repeated with additional replication slaves.

INSERT triggers for BLACKHOLE tables work as expected. However, because the BLACKHOLE table does not actually store any data, UPDATE and DELETE triggers are not activated: The FOR EACH ROW clause in the trigger definition does not apply because there are no rows.

Other possible uses for the BLACKHOLE storage engine include:

- Verification of dump file syntax.
- Measurement of the overhead from binary logging, by comparing performance using BLACKHOLE with and without binary logging enabled.

- `BLACKHOLE` is essentially a “no-op” storage engine, so it could be used for finding performance bottlenecks not related to the storage engine itself.

The `BLACKHOLE` engine is transaction-aware, in the sense that committed transactions are written to the binary log and rolled-back transactions are not.

Blackhole Engine and Auto Increment Columns

The Blackhole engine is a no-op engine. Any operations performed on a table using Blackhole will have no effect. This should be born in mind when considering the behavior of primary key columns that auto increment. The engine will not automatically increment field values, and does not retain auto increment field state. This has important implications in replication.

Consider the following replication scenario where all three of the following conditions apply:

1. On a master server there is a blackhole table with an auto increment field that is a primary key.
2. On a slave the same table exists but using the MyISAM engine.
3. Inserts are performed into the master's table without explicitly setting the auto increment value in the `INSERT` statement itself or through using a `SET INSERT_ID` statement.

In this scenario replication will fail with a duplicate entry error on the primary key column.

In statement based replication, the value of `INSERT_ID` in the context event will always be the same. Replication will therefore fail due to trying insert a row with a duplicate value for a primary key column.

In row based replication, the value that the engine returns for the row always be the same for each insert. This will result in the slave attempting to replay two insert log entries using the same value for the primary key column, and so replication will fail.

Column Filtering

When using row-based replication, (`binlog_format=ROW`), a slave where the last columns are missing from a table is supported, as described in the section Section 18.4.1.10, “Replication with Differing Table Definitions on Master and Slave”.

This filtering works on the slave side, that is, the columns are copied to the slave before they are filtered out. There are at least two cases where it is not desirable to copy the columns to the slave:

1. If the data is confidential, so the slave server should not have access to it.
2. If the master has many slaves, filtering before sending to the slaves may reduce network traffic.

Master column filtering can be achieved using the `BLACKHOLE` engine. This is carried out in a way similar to how master table filtering is achieved - by using the `BLACKHOLE` engine and the `--replicate-do-table` or `--replicate-ignore-table` option.

The setup for the master is:

```
CREATE TABLE t1 (public_col_1, ..., public_col_N,
                 secret_col_1, ..., secret_col_M) ENGINE=MyISAM;
```

The setup for the trusted slave is:

```
CREATE TABLE t1 (public_col_1, ..., public_col_N) ENGINE=BLACKHOLE;
```

The setup for the untrusted slave is:

```
CREATE TABLE t1 (public_col_1, ..., public_col_N) ENGINE=MyISAM;
```

16.7 The MERGE Storage Engine

The MERGE storage engine, also known as the MRG_MyISAM engine, is a collection of identical MyISAM tables that can be used as one. “Identical” means that all tables have identical column and index information. You cannot merge MyISAM tables in which the columns are listed in a different order, do not have exactly the same columns, or have the indexes in different order. However, any or all of the MyISAM tables can be compressed with myisampack. See Section 4.6.6, “myisampack — Generate Compressed, Read-Only MyISAM Tables”. Differences in table options such as AVG_ROW_LENGTH, MAX_ROWS, or PACK_KEYS do not matter.

An alternative to a MERGE table is a partitioned table, which stores partitions of a single table in separate files and enables some operations to be performed more efficiently. For more information, see Chapter 22, *Partitioning*.

When you create a MERGE table, MySQL creates a .MRG file on disk that contains the names of the underlying MyISAM tables that should be used as one. The table format of the MERGE table is stored in the MySQL data dictionary. The underlying tables do not have to be in the same database as the MERGE table.

You can use SELECT, DELETE, UPDATE, and INSERT on MERGE tables. You must have SELECT, DELETE, and UPDATE privileges on the MyISAM tables that you map to a MERGE table.

Note

The use of MERGE tables entails the following security issue: If a user has access to MyISAM table *t*, that user can create a MERGE table *m* that accesses *t*. However, if the user's privileges on *t* are subsequently revoked, the user can continue to access *t* by doing so through *m*.

Use of DROP TABLE with a MERGE table drops only the MERGE specification. The underlying tables are not affected.

To create a MERGE table, you must specify a UNION=(*list-of-tables*) option that indicates which MyISAM tables to use. You can optionally specify an INSERT_METHOD option to control how inserts into the MERGE table take place. Use a value of FIRST or LAST to cause inserts to be made in the first or last underlying table, respectively. If you specify no INSERT_METHOD option or if you specify it with a value of NO, inserts into the MERGE table are not permitted and attempts to do so result in an error.

The following example shows how to create a MERGE table:

```
mysql> CREATE TABLE t1 (
    ->    a INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    ->    message CHAR(20)) ENGINE=MyISAM;
mysql> CREATE TABLE t2 (
    ->    a INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    ->    message CHAR(20)) ENGINE=MyISAM;
mysql> INSERT INTO t1 (message) VALUES ('Testing'),('table'),('t1');
mysql> INSERT INTO t2 (message) VALUES ('Testing'),('table'),('t2');
mysql> CREATE TABLE total (
    ->    a INT NOT NULL AUTO_INCREMENT,
    ->    message CHAR(20), INDEX(a))
    ->    ENGINE=MERGE UNION=(t1,t2) INSERT_METHOD=LAST;
```

Column a is indexed as a PRIMARY KEY in the underlying MyISAM tables, but not in the MERGE table. There it is indexed but not as a PRIMARY KEY because a MERGE table cannot enforce uniqueness over the set of underlying tables. (Similarly, a column with a UNIQUE index in the underlying tables should be indexed in the MERGE table but not as a UNIQUE index.)

After creating the MERGE table, you can use it to issue queries that operate on the group of tables as a whole:

```
mysql> SELECT * FROM total;
+---+---------+
```

```
| a | message |
+---+---------+
| 1 | Testing |
| 2 | table   |
| 3 | t1      |
| 1 | Testing |
| 2 | table   |
| 3 | t2      |
+---+---------+
```

To remap a `MERGE` table to a different collection of `MyISAM` tables, you can use one of the following methods:

- `DROP` the `MERGE` table and re-create it.
- Use `ALTER TABLE tbl_name UNION=(...)` to change the list of underlying tables.

  It is also possible to use `ALTER TABLE ... UNION=()` (that is, with an empty `UNION` clause) to remove all of the underlying tables. However, in this case, the table is effectively empty and inserts fail because there is no underlying table to take new rows. Such a table might be useful as a template for creating new `MERGE` tables with `CREATE TABLE ... LIKE`.

The underlying table definitions and indexes must conform closely to the definition of the `MERGE` table. Conformance is checked when a table that is part of a `MERGE` table is opened, not when the `MERGE` table is created. If any table fails the conformance checks, the operation that triggered the opening of the table fails. This means that changes to the definitions of tables within a `MERGE` may cause a failure when the `MERGE` table is accessed. The conformance checks applied to each table are:

- The underlying table and the `MERGE` table must have the same number of columns.
- The column order in the underlying table and the `MERGE` table must match.
- Additionally, the specification for each corresponding column in the parent `MERGE` table and the underlying tables are compared and must satisfy these checks:
  - The column type in the underlying table and the `MERGE` table must be equal.
  - The column length in the underlying table and the `MERGE` table must be equal.
  - The column of the underlying table and the `MERGE` table can be `NULL`.
- The underlying table must have at least as many indexes as the `MERGE` table. The underlying table may have more indexes than the `MERGE` table, but cannot have fewer.

  Note

  A known issue exists where indexes on the same columns must be in identical order, in both the `MERGE` table and the underlying `MyISAM` table. See Bug #33653.

  Each index must satisfy these checks:

  - The index type of the underlying table and the `MERGE` table must be the same.
  - The number of index parts (that is, multiple columns within a compound index) in the index definition for the underlying table and the `MERGE` table must be the same.
  - For each index part:
    - Index part lengths must be equal.
    - Index part types must be equal.
    - Index part languages must be equal.

- Check whether index parts can be `NULL`.

If a `MERGE` table cannot be opened or used because of a problem with an underlying table, `CHECK TABLE` displays information about which table caused the problem.

Additional Resources

- A forum dedicated to the `MERGE` storage engine is available at http://forums.mysql.com/list.php?93.

16.7.1 MERGE Table Advantages and Disadvantages

`MERGE` tables can help you solve the following problems:

- Easily manage a set of log tables. For example, you can put data from different months into separate tables, compress some of them with `myisampack`, and then create a `MERGE` table to use them as one.
- Obtain more speed. You can split a large read-only table based on some criteria, and then put individual tables on different disks. A `MERGE` table structured this way could be much faster than using a single large table.
- Perform more efficient searches. If you know exactly what you are looking for, you can search in just one of the underlying tables for some queries and use a `MERGE` table for others. You can even have many different `MERGE` tables that use overlapping sets of tables.
- Perform more efficient repairs. It is easier to repair individual smaller tables that are mapped to a `MERGE` table than to repair a single large table.
- Instantly map many tables as one. A `MERGE` table need not maintain an index of its own because it uses the indexes of the individual tables. As a result, `MERGE` table collections are *very* fast to create or remap. (You must still specify the index definitions when you create a `MERGE` table, even though no indexes are created.)
- If you have a set of tables from which you create a large table on demand, you can instead create a `MERGE` table from them on demand. This is much faster and saves a lot of disk space.
- Exceed the file size limit for the operating system. Each `MyISAM` table is bound by this limit, but a collection of `MyISAM` tables is not.
- You can create an alias or synonym for a `MyISAM` table by defining a `MERGE` table that maps to that single table. There should be no really notable performance impact from doing this (only a couple of indirect calls and `memcpy()` calls for each read).

The disadvantages of `MERGE` tables are:

- You can use only identical `MyISAM` tables for a `MERGE` table.
- Some `MyISAM` features are unavailable in `MERGE` tables. For example, you cannot create `FULLTEXT` indexes on `MERGE` tables. (You can create `FULLTEXT` indexes on the underlying `MyISAM` tables, but you cannot search the `MERGE` table with a full-text search.)
- If the `MERGE` table is nontemporary, all underlying `MyISAM` tables must be nontemporary. If the `MERGE` table is temporary, the `MyISAM` tables can be any mix of temporary and nontemporary.
- `MERGE` tables use more file descriptors than `MyISAM` tables. If 10 clients are using a `MERGE` table that maps to 10 tables, the server uses (10 × 10) + 10 file descriptors. (10 data file descriptors for each of the 10 clients, and 10 index file descriptors shared among the clients.)
- Index reads are slower. When you read an index, the `MERGE` storage engine needs to issue a read on all underlying tables to check which one most closely matches a given index value. To read

the next index value, the MERGE storage engine needs to search the read buffers to find the next value. Only when one index buffer is used up does the storage engine need to read the next index block. This makes MERGE indexes much slower on eq_ref searches, but not much slower on ref searches. For more information about eq_ref and ref, see Section 13.8.2, “EXPLAIN Syntax”.

16.7.2 MERGE Table Problems

The following are known problems with MERGE tables:

- In versions of MySQL Server prior to 5.1.23, it was possible to create temporary merge tables with nontemporary child MyISAM tables.

  From versions 5.1.23, MERGE children were locked through the parent table. If the parent was temporary, it was not locked and so the children were not locked either. Parallel use of the MyISAM tables corrupted them.

- If you use ALTER TABLE to change a MERGE table to another storage engine, the mapping to the underlying tables is lost. Instead, the rows from the underlying MyISAM tables are copied into the altered table, which then uses the specified storage engine.

- The INSERT_METHOD table option for a MERGE table indicates which underlying MyISAM table to use for inserts into the MERGE table. However, use of the AUTO_INCREMENT table option for that MyISAM table has no effect for inserts into the MERGE table until at least one row has been inserted directly into the MyISAM table.

- A MERGE table cannot maintain uniqueness constraints over the entire table. When you perform an INSERT, the data goes into the first or last MyISAM table (as determined by the INSERT_METHOD option). MySQL ensures that unique key values remain unique within that MyISAM table, but not over all the underlying tables in the collection.

- Because the MERGE engine cannot enforce uniqueness over the set of underlying tables, REPLACE does not work as expected. The two key facts are:

  - REPLACE can detect unique key violations only in the underlying table to which it is going to write (which is determined by the INSERT_METHOD option). This differs from violations in the MERGE table itself.

  - If REPLACE detects a unique key violation, it will change only the corresponding row in the underlying table it is writing to; that is, the first or last table, as determined by the INSERT_METHOD option.

  Similar considerations apply for INSERT ... ON DUPLICATE KEY UPDATE.

- MERGE tables do not support partitioning. That is, you cannot partition a MERGE table, nor can any of a MERGE table's underlying MyISAM tables be partitioned.

- You should not use ANALYZE TABLE, REPAIR TABLE, OPTIMIZE TABLE, ALTER TABLE, DROP TABLE, DELETE without a WHERE clause, or TRUNCATE TABLE on any of the tables that are mapped into an open MERGE table. If you do so, the MERGE table may still refer to the original table and yield unexpected results. To work around this problem, ensure that no MERGE tables remain open by issuing a FLUSH TABLES statement prior to performing any of the named operations.

  The unexpected results include the possibility that the operation on the MERGE table will report table corruption. If this occurs after one of the named operations on the underlying MyISAM tables, the corruption message is spurious. To deal with this, issue a FLUSH TABLES statement after modifying the MyISAM tables.

- DROP TABLE on a table that is in use by a MERGE table does not work on Windows because the MERGE storage engine's table mapping is hidden from the upper layer of MySQL. Windows does not permit open files to be deleted, so you first must flush all MERGE tables (with FLUSH TABLES) or drop the MERGE table before dropping the table.

- The definition of the MyISAM tables and the MERGE table are checked when the tables are accessed (for example, as part of a SELECT or INSERT statement). The checks ensure that the definitions of the tables and the parent MERGE table definition match by comparing column order, types, sizes and associated indexes. If there is a difference between the tables, an error is returned and the statement fails. Because these checks take place when the tables are opened, any changes to the definition of a single table, including column changes, column ordering, and engine alterations will cause the statement to fail.
- The order of indexes in the MERGE table and its underlying tables should be the same. If you use ALTER TABLE to add a UNIQUE index to a table used in a MERGE table, and then use ALTER TABLE to add a nonunique index on the MERGE table, the index ordering is different for the tables if there was already a nonunique index in the underlying table. (This happens because ALTER TABLE puts UNIQUE indexes before nonunique indexes to facilitate rapid detection of duplicate keys.) Consequently, queries on tables with such indexes may return unexpected results.
- If you encounter an error message similar to ERROR 1017 (HY000): Can't find file: '*tbl_name*.MRG' (errno: 2), it generally indicates that some of the underlying tables do not use the MyISAM storage engine. Confirm that all of these tables are MyISAM.
- The maximum number of rows in a MERGE table is 2^{64} (~1.844E+19; the same as for a MyISAM table). It is not possible to merge multiple MyISAM tables into a single MERGE table that would have more than this number of rows.
- Use of underlying MyISAM tables of differing row formats with a parent MERGE table is currently known to fail. See Bug #32364.
- You cannot change the union list of a nontemporary MERGE table when LOCK TABLES is in effect. The following does *not* work:

  ```
  CREATE TABLE m1 ... ENGINE=MRG_MYISAM ...;
  LOCK TABLES t1 WRITE, t2 WRITE, m1 WRITE;
  ALTER TABLE m1 ... UNION=(t1,t2) ...;
  ```

  However, you can do this with a temporary MERGE table.
- You cannot create a MERGE table with CREATE ... SELECT, neither as a temporary MERGE table, nor as a nontemporary MERGE table. For example:

  ```
  CREATE TABLE m1 ... ENGINE=MRG_MYISAM ... SELECT ...;
  ```

  Attempts to do this result in an error: *tbl_name* is not BASE TABLE.
- In some cases, differing PACK_KEYS table option values among the MERGE and underlying tables cause unexpected results if the underlying tables contain CHAR or BINARY columns. As a workaround, use ALTER TABLE to ensure that all involved tables have the same PACK_KEYS value. (Bug #50646)

16.8 The FEDERATED Storage Engine

The FEDERATED storage engine lets you access data from a remote MySQL database without using replication or cluster technology. Querying a local FEDERATED table automatically pulls the data from the remote (federated) tables. No data is stored on the local tables.

To include the FEDERATED storage engine if you build MySQL from source, invoke CMake with the -DWITH_FEDERATED_STORAGE_ENGINE option.

The FEDERATED storage engine is not enabled by default in the running server; to enable FEDERATED, you must start the MySQL server binary using the --federated option.

To examine the source for the FEDERATED engine, look in the storage/federated directory of a MySQL source distribution.

16.8.1 FEDERATED Storage Engine Overview

When you create a table using one of the standard storage engines (such as `MyISAM`, `CSV` or `InnoDB`), the table consists of the table definition and the associated data. When you create a `FEDERATED` table, the table definition is the same, but the physical storage of the data is handled on a remote server.

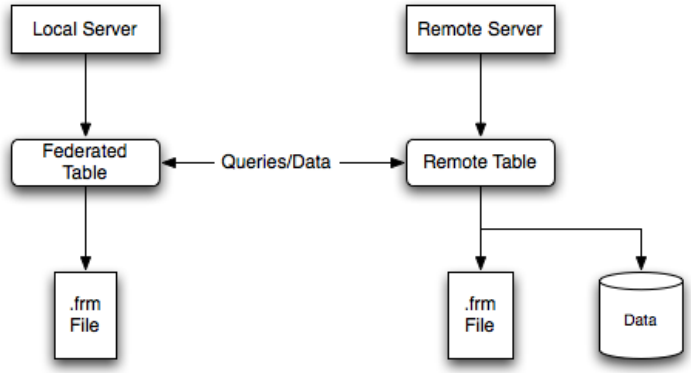
A `FEDERATED` table consists of two elements:

- A *remote server* with a database table, which in turn consists of the table definition (stored in the MySQL data dictionary) and the associated table. The table type of the remote table may be any type supported by the remote `mysqld` server, including `MyISAM` or `InnoDB`.
- A *local server* with a database table, where the table definition matches that of the corresponding table on the remote server. The table definition is stored in the data dictionary. There is no data file on the local server. Instead, the table definition includes a connection string that points to the remote table.

When executing queries and statements on a `FEDERATED` table on the local server, the operations that would normally insert, update or delete information from a local data file are instead sent to the remote server for execution, where they update the data file on the remote server or return matching rows from the remote server.

The basic structure of a `FEDERATED` table setup is shown in Figure 16.1, “FEDERATED Table Structure”.

Figure 16.1 FEDERATED Table Structure



When a client issues an SQL statement that refers to a `FEDERATED` table, the flow of information between the local server (where the SQL statement is executed) and the remote server (where the data is physically stored) is as follows:

1. The storage engine looks through each column that the `FEDERATED` table has and constructs an appropriate SQL statement that refers to the remote table.
2. The statement is sent to the remote server using the MySQL client API.
3. The remote server processes the statement and the local server retrieves any result that the statement produces (an affected-rows count or a result set).
4. If the statement produces a result set, each column is converted to internal storage engine format that the `FEDERATED` engine expects and can use to display the result to the client that issued the original statement.

The local server communicates with the remote server using MySQL client C API functions. It invokes `mysql_real_query()` to send the statement. To read a result set, it uses `mysql_store_result()` and fetches rows one at a time using `mysql_fetch_row()`.

16.8.2 How to Create FEDERATED Tables

To create a `FEDERATED` table you should follow these steps:

1. Create the table on the remote server. Alternatively, make a note of the table definition of an existing table, perhaps using the `SHOW CREATE TABLE` statement.

2. Create the table on the local server with an identical table definition, but adding the connection information that links the local table to the remote table.

For example, you could create the following table on the remote server:

```
CREATE TABLE test_table (
    id     INT(20) NOT NULL AUTO_INCREMENT,
    name   VARCHAR(32) NOT NULL DEFAULT '',
    other  INT(20) NOT NULL DEFAULT '0',
    PRIMARY KEY  (id),
    INDEX name (name),
    INDEX other_key (other)
)
ENGINE=MyISAM
DEFAULT CHARSET=utf8mb4;
```

To create the local table that will be federated to the remote table, there are two options available. You can either create the local table and specify the connection string (containing the server name, login, password) to be used to connect to the remote table using the `CONNECTION`, or you can use an existing connection that you have previously created using the `CREATE SERVER` statement.

Important

When you create the local table it *must* have an identical field definition to the remote table.

Note

You can improve the performance of a `FEDERATED` table by adding indexes to the table on the host. The optimization will occur because the query sent to the remote server will include the contents of the `WHERE` clause and will be sent to the remote server and subsequently executed locally. This reduces the network traffic that would otherwise request the entire table from the server for local processing.

16.8.2.1 Creating a FEDERATED Table Using CONNECTION

To use the first method, you must specify the `CONNECTION` string after the engine type in a `CREATE TABLE` statement. For example:

```
CREATE TABLE federated_table (
    id     INT(20) NOT NULL AUTO_INCREMENT,
    name   VARCHAR(32) NOT NULL DEFAULT '',
    other  INT(20) NOT NULL DEFAULT '0',
    PRIMARY KEY  (id),
    INDEX name (name),
    INDEX other_key (other)
)
ENGINE=FEDERATED
DEFAULT CHARSET=utf8mb4
CONNECTION='mysql://fed_user@remote_host:9306/federated/test_table';
```

Note

`CONNECTION` replaces the `COMMENT` used in some previous versions of MySQL.

The CONNECTION string contains the information required to connect to the remote server containing the table that will be used to physically store the data. The connection string specifies the server name, login credentials, port number and database/table information. In the example, the remote table is on the server remote_host, using port 9306. The name and port number should match the host name (or IP address) and port number of the remote MySQL server instance you want to use as your remote table.

The format of the connection string is as follows:

```
scheme://user_name[:password]@host_name[:port_num]/db_name/tbl_name
```

Where:

- *scheme*: A recognized connection protocol. Only mysql is supported as the *scheme* value at this point.
- *user_name*: The user name for the connection. This user must have been created on the remote server, and must have suitable privileges to perform the required actions (SELECT, INSERT, UPDATE, and so forth) on the remote table.
- *password*: (Optional) The corresponding password for *user_name*.
- *host_name*: The host name or IP address of the remote server.
- *port_num*: (Optional) The port number for the remote server. The default is 3306.
- *db_name*: The name of the database holding the remote table.
- *tbl_name*: The name of the remote table. The name of the local and the remote table do not have to match.

Sample connection strings:

```
CONNECTION='mysql://username:password@hostname:port/database/tablename'
CONNECTION='mysql://username@hostname/database/tablename'
CONNECTION='mysql://username:password@hostname/database/tablename'
```

16.8.2.2 Creating a FEDERATED Table Using CREATE SERVER

If you are creating a number of FEDERATED tables on the same server, or if you want to simplify the process of creating FEDERATED tables, you can use the CREATE SERVER statement to define the server connection parameters, just as you would with the CONNECTION string.

The format of the CREATE SERVER statement is:

```
CREATE SERVER
server_name
FOREIGN DATA WRAPPER wrapper_name
OPTIONS (option [, option] ...)
```

The *server_name* is used in the connection string when creating a new FEDERATED table.

For example, to create a server connection identical to the CONNECTION string:

```
CONNECTION='mysql://fed_user@remote_host:9306/federated/test_table';
```

You would use the following statement:

```
CREATE SERVER fedlink
FOREIGN DATA WRAPPER mysql
OPTIONS (USER 'fed_user', HOST 'remote_host', PORT 9306, DATABASE 'federated');
```

To create a FEDERATED table that uses this connection, you still use the CONNECTION keyword, but specify the name you used in the CREATE SERVER statement.

```
CREATE TABLE test_table (
    id     INT(20) NOT NULL AUTO_INCREMENT,
    name   VARCHAR(32) NOT NULL DEFAULT '',
    other  INT(20) NOT NULL DEFAULT '0',
    PRIMARY KEY  (id),
    INDEX name (name),
    INDEX other_key (other)
)
ENGINE=FEDERATED
DEFAULT CHARSET=utf8mb4
CONNECTION='fedlink/test_table';
```

The connection name in this example contains the name of the connection (`fedlink`) and the name of the table (`test_table`) to link to, separated by a slash. If you specify only the connection name without a table name, the table name of the local table is used instead.

For more information on `CREATE SERVER`, see Section 13.1.14, “CREATE SERVER Syntax”.

The `CREATE SERVER` statement accepts the same arguments as the `CONNECTION` string. The `CREATE SERVER` statement updates the rows in the `mysql.servers` table. See the following table for information on the correspondence between parameters in a connection string, options in the `CREATE SERVER` statement, and the columns in the `mysql.servers` table. For reference, the format of the `CONNECTION` string is as follows:

```
scheme://user_name[:password]@host_name[:port_num]/db_name/tbl_name
```

Description	`CONNECTION` string	`CREATE SERVER` option	`mysql.servers` column
Connection scheme	*`scheme`*	`wrapper_name`	`Wrapper`
Remote user	*`user_name`*	`USER`	`Username`
Remote password	*`password`*	`PASSWORD`	`Password`
Remote host	*`host_name`*	`HOST`	`Host`
Remote port	*`port_num`*	`PORT`	`Port`
Remote database	*`db_name`*	`DATABASE`	`Db`

16.8.3 FEDERATED Storage Engine Notes and Tips

You should be aware of the following points when using the `FEDERATED` storage engine:

- `FEDERATED` tables may be replicated to other slaves, but you must ensure that the slave servers are able to use the user/password combination that is defined in the `CONNECTION` string (or the row in the `mysql.servers` table) to connect to the remote server.

The following items indicate features that the `FEDERATED` storage engine does and does not support:

- The remote server must be a MySQL server.
- The remote table that a `FEDERATED` table points to *must* exist before you try to access the table through the `FEDERATED` table.
- It is possible for one `FEDERATED` table to point to another, but you must be careful not to create a loop.
- A `FEDERATED` table does not support indexes in the usual sense; because access to the table data is handled remotely, it is actually the remote table that makes use of indexes. This means that, for a query that cannot use any indexes and so requires a full table scan, the server fetches all rows from the remote table and filters them locally. This occurs regardless of any `WHERE` or `LIMIT` used with this `SELECT` statement; these clauses are applied locally to the returned rows.

Queries that fail to use indexes can thus cause poor performance and network overload. In addition, since returned rows must be stored in memory, such a query can also lead to the local server swapping, or even hanging.

- Care should be taken when creating a `FEDERATED` table since the index definition from an equivalent `MyISAM` or other table may not be supported. For example, creating a `FEDERATED` table with an index prefix on `VARCHAR`, `TEXT` or `BLOB` columns will fail. The following definition in `MyISAM` is valid:

```
CREATE TABLE `T1`(`A` VARCHAR(100),UNIQUE KEY(`A`(30))) ENGINE=MYISAM;
```

The key prefix in this example is incompatible with the `FEDERATED` engine, and the equivalent statement will fail:

```
CREATE TABLE `T1`(`A` VARCHAR(100),UNIQUE KEY(`A`(30))) ENGINE=FEDERATED
  CONNECTION='MYSQL://127.0.0.1:3306/TEST/T1';
```

If possible, you should try to separate the column and index definition when creating tables on both the remote server and the local server to avoid these index issues.

- Internally, the implementation uses `SELECT`, `INSERT`, `UPDATE`, and `DELETE`, but not `HANDLER`.
- The `FEDERATED` storage engine supports `SELECT`, `INSERT`, `UPDATE`, `DELETE`, `TRUNCATE TABLE`, and indexes. It does not support `ALTER TABLE`, or any Data Definition Language statements that directly affect the structure of the table, other than `DROP TABLE`. The current implementation does not use prepared statements.
- `FEDERATED` accepts `INSERT ... ON DUPLICATE KEY UPDATE` statements, but if a duplicate-key violation occurs, the statement fails with an error.
- Performance on a `FEDERATED` table when performing bulk inserts (for example, on a `INSERT INTO ... SELECT ...` statement) is slower than with other table types because each selected row is treated as an individual `INSERT` statement on the `FEDERATED` table.
- Transactions are not supported.
- `FEDERATED` performs bulk-insert handling such that multiple rows are sent to the remote table in a batch. This provides a performance improvement and enables the remote table to perform improvement. Also, if the remote table is transactional, it enables the remote storage engine to perform statement rollback properly should an error occur. This capability has the following limitations:
  - The size of the insert cannot exceed the maximum packet size between servers. If the insert exceeds this size, it is broken into multiple packets and the rollback problem can occur.
  - Bulk-insert handling does not occur for `INSERT ... ON DUPLICATE KEY UPDATE`.
- There is no way for the `FEDERATED` engine to know if the remote table has changed. The reason for this is that this table must work like a data file that would never be written to by anything other than the database system. The integrity of the data in the local table could be breached if there was any change to the remote database.
- When using a `CONNECTION` string, you cannot use an '@' character in the password. You can get round this limitation by using the `CREATE SERVER` statement to create a server connection.
- The `insert_id` and `timestamp` options are not propagated to the data provider.
- Any `DROP TABLE` statement issued against a `FEDERATED` table drops only the local table, not the remote table.
- `FEDERATED` tables do not work with the query cache.
- User-defined partitioning is not supported for `FEDERATED` tables.

16.8.4 FEDERATED Storage Engine Resources

The following additional resources are available for the FEDERATED storage engine:

- A forum dedicated to the FEDERATED storage engine is available at http://forums.mysql.com/list.php?105.

16.9 The EXAMPLE Storage Engine

The EXAMPLE storage engine is a stub engine that does nothing. Its purpose is to serve as an example in the MySQL source code that illustrates how to begin writing new storage engines. As such, it is primarily of interest to developers.

To enable the EXAMPLE storage engine if you build MySQL from source, invoke CMake with the -DWITH_EXAMPLE_STORAGE_ENGINE option.

To examine the source for the EXAMPLE engine, look in the storage/example directory of a MySQL source distribution.

When you create an EXAMPLE table, no files are created. No data can be stored into the table. Retrievals return an empty result.

```
mysql> CREATE TABLE test (i INT) ENGINE = EXAMPLE;
Query OK, 0 rows affected (0.78 sec)

mysql> INSERT INTO test VALUES(1),(2),(3);
ERROR 1031 (HY000): Table storage engine for 'test' doesn't »
                    have this option

mysql> SELECT * FROM test;
Empty set (0.31 sec)
```

The EXAMPLE storage engine does not support indexing.

The EXAMPLE storage engine does not support partitioning.

16.10 Other Storage Engines

Other storage engines may be available from third parties and community members that have used the Custom Storage Engine interface.

Third party engines are not supported by MySQL. For further information, documentation, installation guides, bug reporting or for any help or assistance with these engines, please contact the developer of the engine directly.

For more information on developing a customer storage engine that can be used with the Pluggable Storage Engine Architecture, see MySQL Internals: Writing a Custom Storage Engine.

16.11 Overview of MySQL Storage Engine Architecture

The MySQL pluggable storage engine architecture enables a database professional to select a specialized storage engine for a particular application need while being completely shielded from the need to manage any specific application coding requirements. The MySQL server architecture isolates the application programmer and DBA from all of the low-level implementation details at the storage level, providing a consistent and easy application model and API. Thus, although there are different capabilities across different storage engines, the application is shielded from these differences.

The pluggable storage engine architecture provides a standard set of management and support services that are common among all underlying storage engines. The storage engines themselves

are the components of the database server that actually perform actions on the underlying data that is maintained at the physical server level.

This efficient and modular architecture provides huge benefits for those wishing to specifically target a particular application need—such as data warehousing, transaction processing, or high availability situations—while enjoying the advantage of utilizing a set of interfaces and services that are independent of any one storage engine.

The application programmer and DBA interact with the MySQL database through Connector APIs and service layers that are above the storage engines. If application changes bring about requirements that demand the underlying storage engine change, or that one or more storage engines be added to support new needs, no significant coding or process changes are required to make things work. The MySQL server architecture shields the application from the underlying complexity of the storage engine by presenting a consistent and easy-to-use API that applies across storage engines.

16.11.1 Pluggable Storage Engine Architecture

MySQL Server uses a pluggable storage engine architecture that enables storage engines to be loaded into and unloaded from a running MySQL server.

Plugging in a Storage Engine

Before a storage engine can be used, the storage engine plugin shared library must be loaded into MySQL using the `INSTALL PLUGIN` statement. For example, if the `EXAMPLE` engine plugin is named `example` and the shared library is named `ha_example.so`, you load it with the following statement:

```
mysql> INSTALL PLUGIN example SONAME 'ha_example.so';
```

To install a pluggable storage engine, the plugin file must be located in the MySQL plugin directory, and the user issuing the `INSTALL PLUGIN` statement must have `INSERT` privilege for the `mysql.plugin` table.

The shared library must be located in the MySQL server plugin directory, the location of which is given by the `plugin_dir` system variable.

Unplugging a Storage Engine

To unplug a storage engine, use the `UNINSTALL PLUGIN` statement:

```
mysql> UNINSTALL PLUGIN example;
```

If you unplug a storage engine that is needed by existing tables, those tables become inaccessible, but will still be present on disk (where applicable). Ensure that there are no tables using a storage engine before you unplug the storage engine.

16.11.2 The Common Database Server Layer

A MySQL pluggable storage engine is the component in the MySQL database server that is responsible for performing the actual data I/O operations for a database as well as enabling and enforcing certain feature sets that target a specific application need. A major benefit of using specific storage engines is that you are only delivered the features needed for a particular application, and therefore you have less system overhead in the database, with the end result being more efficient and higher database performance. This is one of the reasons that MySQL has always been known to have such high performance, matching or beating proprietary monolithic databases in industry standard benchmarks.

From a technical perspective, what are some of the unique supporting infrastructure components that are in a storage engine? Some of the key feature differentiations include:

- *Concurrency*: Some applications have more granular lock requirements (such as row-level locks) than others. Choosing the right locking strategy can reduce overhead and therefore improve overall performance. This area also includes support for capabilities such as multi-version concurrency control or “snapshot” read.
- *Transaction Support*: Not every application needs transactions, but for those that do, there are very well defined requirements such as ACID compliance and more.
- *Referential Integrity*: The need to have the server enforce relational database referential integrity through DDL defined foreign keys.
- *Physical Storage*: This involves everything from the overall page size for tables and indexes as well as the format used for storing data to physical disk.
- *Index Support*: Different application scenarios tend to benefit from different index strategies. Each storage engine generally has its own indexing methods, although some (such as B-tree indexes) are common to nearly all engines.
- *Memory Caches*: Different applications respond better to some memory caching strategies than others, so although some memory caches are common to all storage engines (such as those used for user connections or MySQL's high-speed Query Cache), others are uniquely defined only when a particular storage engine is put in play.
- *Performance Aids*: This includes multiple I/O threads for parallel operations, thread concurrency, database checkpointing, bulk insert handling, and more.
- *Miscellaneous Target Features*: This may include support for geospatial operations, security restrictions for certain data manipulation operations, and other similar features.

Each set of the pluggable storage engine infrastructure components are designed to offer a selective set of benefits for a particular application. Conversely, avoiding a set of component features helps reduce unnecessary overhead. It stands to reason that understanding a particular application's set of requirements and selecting the proper MySQL storage engine can have a dramatic impact on overall system efficiency and performance.

Chapter 17 High Availability and Scalability

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Data is the currency of today's web, mobile, social, enterprise and cloud applications. Ensuring data is always available is a top priority for any organization. Minutes of downtime can result in significant loss of revenue and reputation.

There is no “one size fits all” approach to delivering High Availability (HA). Unique application attributes, business requirements, operational capabilities and legacy infrastructure can all influence HA technology selection. And technology is only one element in delivering HA: people and processes are just as critical as the technology itself.

MySQL is deployed into many applications demanding availability and scalability. **Availability** refers to the ability to cope with, and if necessary recover from, failures on the host, including failures of MySQL, the operating system, or the hardware and maintenance activity that may otherwise cause downtime. Scalability refers to the ability to spread both the database and the load of your application queries across multiple MySQL servers.

Because each application has different operational and availability requirements, MySQL offers a range of certified and supported solutions, delivering the appropriate levels of High Availability (HA) and scalability to meet service level requirements. Such solutions extend from replication, through virtualization and geographically redundant, multi-data center solutions delivering 99.999% uptime.

Selecting the right high availability solution for an application largely depends on:

- The level of availability required.
- The type of application being deployed.
- Accepted best practices within your own environment.

The primary solutions supported by MySQL include:

- MySQL Replication. Learn more: Chapter 18, *Replication*.
- MySQL Fabric. Learn more: MySQL Fabric.
- MySQL Cluster. Learn more: MySQL NDB Cluster 7.5 and NDB Cluster 7.6.
- Oracle MySQL Cloud Service. Learn more about MySQL Cloud Service.
- Oracle Clusterware Agent for MySQL. Learn more about Oracle Clusterware.
- MySQL with Solaris Cluster. Learn more about Solaris Cluster.

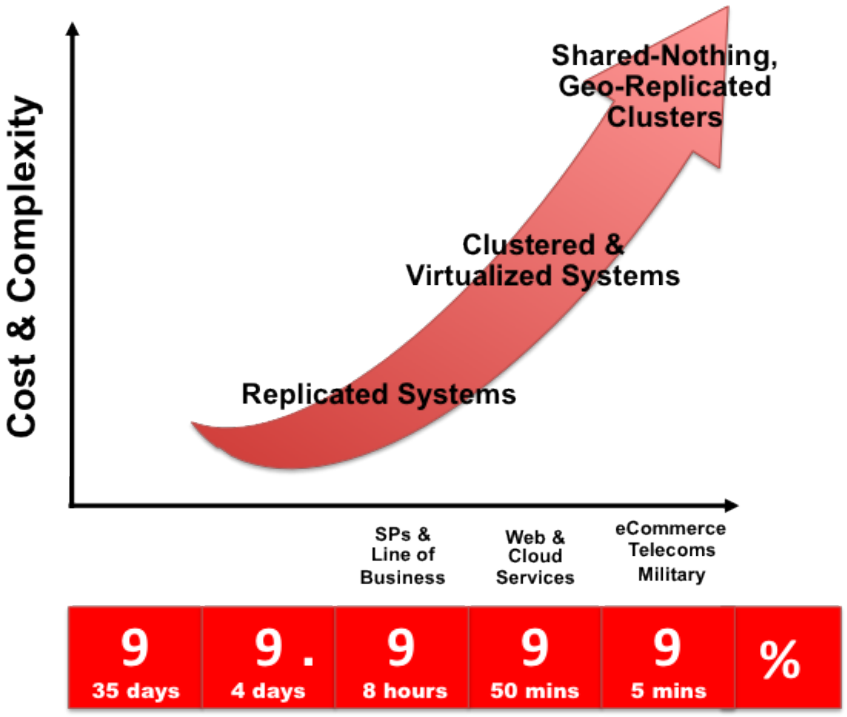
Further options are available using third-party solutions.

Each architecture used to achieve highly available database services is differentiated by the levels of uptime it offers. These architectures can be grouped into three main categories:

- Data Replication.
- Clustered & Virtualized Systems.
- Shared-Nothing, Geographically-Replicated Clusters.

As illustrated in the following figure, each of these architectures offers progressively higher levels of uptime, which must be balanced against potentially greater levels of cost and complexity that each can incur. Simply deploying a high availability architecture is not a guarantee of actually delivering HA. In fact, a poorly implemented and maintained shared-nothing cluster could easily deliver lower levels of availability than a simple data replication solution.

Figure 17.1 Tradeoffs: Cost and Complexity versus Availability



The following table compares the HA and Scalability capabilities of the various MySQL solutions:

Table 17.1 Feature Comparison of MySQL HA Solutions

Requirement	MySQL Replication	MySQL Cluster
Availability		
Platform Support	All Supported by MySQL Server (http://www.mysql.com/support/supportedplatforms/database.html)	All Supported by MySQL Cluster (http://www.mysql.com/support/supportedplatforms/cluster.html)
Automated IP Failover	No	Depends on Connector and Configuration
Automated Database Failover	No	Yes
Automatic Data Resynchronization	No	Yes
Typical Failover Time	User / Script Dependent	1 Second and Less
Synchronous Replication	No, Asynchronous and Semisynchronous	Yes
Shared Storage	No, Distributed	No, Distributed

Requirement	MySQL Replication	MySQL Cluster
Geographic redundancy support	Yes	Yes, via MySQL Replication
Update Schema On-Line	No	Yes
Scalability		
Number of Nodes	One Master, Multiple Slaves	255
Built-in Load Balancing	Reads, via MySQL Replication	Yes, Reads and Writes
Supports Read-Intensive Workloads	Yes	Yes
Supports Write-Intensive Workloads	Yes, via Application-Level Sharding	Yes, via Auto-Sharding
Scale On-Line (add nodes, repartition, etc.)	No	Yes

17.1 Using ZFS Replication

To support high availability environments, providing an instant copy of the information on both the currently active machine and the hot backup is a critical part of the HA solution. There are many solutions to this problem, such as Chapter 18, *Replication*.

The ZFS file system provides functionality to create a snapshot of the file system contents, transfer the snapshot to another machine, and extract the snapshot to recreate the file system. You can create a snapshot at any time, and you can create as many snapshots as you like. By continually creating, transferring, and restoring snapshots, you can provide synchronization between one or more machines in a fashion similar to DRBD.

The following example shows a simple Solaris system running with a single ZFS pool, mounted at `/scratchpool`:

```
Filesystem             size   used  avail capacity  Mounted on
/dev/dsk/c0d0s0        4.6G   3.7G   886M    82%    /
/devices                 0K     0K     0K     0%    /devices
ctfs                     0K     0K     0K     0%    /system/contract
proc                     0K     0K     0K     0%    /proc
mnttab                   0K     0K     0K     0%    /etc/mnttab
swap                   1.4G   892K   1.4G     1%    /etc/svc/volatile
objfs                    0K     0K     0K     0%    /system/object
/usr/lib/libc/libc_hwcap1.so.1
                       4.6G   3.7G   886M    82%    /lib/libc.so.1
fd                       0K     0K     0K     0%    /dev/fd
swap                   1.4G    40K   1.4G     1%    /tmp
swap                   1.4G    28K   1.4G     1%    /var/run
/dev/dsk/c0d0s7         26G   913M    25G     4%    /export/home
scratchpool             16G    24K    16G     1%    /scratchpool
```

The MySQL data is stored in a directory on `/scratchpool`. To help demonstrate some of the basic replication functionality, there are also other items stored in `/scratchpool` as well:

```
total 17
drwxr-xr-x  31 root     bin           50 Jul 21 07:32 DTT/
drwxr-xr-x   4 root     bin            5 Jul 21 07:32 SUNWmlib/
drwxr-xr-x  14 root     sys           16 Nov  5 09:56 SUNWspro/
drwxrwxrwx  19 1000     1000          40 Nov  6 19:16 emacs-22.1/
```

To create a snapshot of the file system, you use `zfs snapshot`, specifying the pool and the snapshot name:

```
root-shell> zfs snapshot scratchpool@snap1
```

To list the snapshots already taken:

```
root-shell> zfs list -t snapshot
NAME                 USED  AVAIL  REFER  MOUNTPOINT
scratchpool@snap1      0      -  24.5K  -
scratchpool@snap2      0      -  24.5K  -
```

The snapshots themselves are stored within the file system metadata, and the space required to keep them varies as time goes on because of the way the snapshots are created. The initial creation of a snapshot is very quick, because instead of taking an entire copy of the data and metadata required to hold the entire snapshot, ZFS records only the point in time and metadata of when the snapshot was created.

As more changes to the original file system are made, the size of the snapshot increases because more space is required to keep the record of the old blocks. If you create lots of snapshots, say one per day, and then delete the snapshots from earlier in the week, the size of the newer snapshots might also increase, as the changes that make up the newer state have to be included in the more recent snapshots, rather than being spread over the seven snapshots that make up the week.

You cannot directly back up the snapshots because they exist within the file system metadata rather than as regular files. To get the snapshot into a format that you can copy to another file system, tape, and so on, you use the `zfs send` command to create a stream version of the snapshot.

For example, to write the snapshot out to a file:

```
root-shell> zfs send scratchpool@snap1 >/backup/scratchpool-snap1
```

Or tape:

```
root-shell> zfs send scratchpool@snap1 >/dev/rmt/0
```

You can also write out the incremental changes between two snapshots using `zfs send`:

```
root-shell> zfs send scratchpool@snap1 scratchpool@snap2 >/backup/scratchpool-changes
```

To recover a snapshot, you use `zfs recv`, which applies the snapshot information either to a new file system, or to an existing one.

17.1.1 Using ZFS for File System Replication

Because `zfs send` and `zfs recv` use streams to exchange data, you can use them to replicate information from one system to another by combining `zfs send`, `ssh`, and `zfs recv`.

For example, to copy a snapshot of the `scratchpool` file system to a new file system called `slavepool` on a new server, you would use the following command. This sequence combines the snapshot of `scratchpool`, the transmission to the slave machine (using `ssh` with login credentials), and the recovery of the snapshot on the slave using `zfs recv`:

```
root-shell> zfs send scratchpool@snap1 |ssh id@host pfexec zfs recv -F slavepool
```

The first part of the pipeline, `zfs send scratchpool@snap1`, streams the snapshot. The `ssh` command, and the command that it executes on the other server, `pfexec zfs recv -F slavepool`, receives the streamed snapshot data and writes it to slavepool. In this instance, I've specified the `-F` option which forces the snapshot data to be applied, and is therefore destructive. This is fine, as I'm creating the first version of my replicated file system.

On the slave machine, the replicated file system contains the exact same content:

```
root-shell> ls -al /slavepool/
total 23
drwxr-xr-x   6 root     root           7 Nov  8 09:13 ./
drwxr-xr-x  29 root     root          34 Nov  9 07:06 ../
drwxr-xr-x  31 root     bin           50 Jul 21 07:32 DTT/
drwxr-xr-x   4 root     bin            5 Jul 21 07:32 SUNWmlib/
drwxr-xr-x  14 root     sys           16 Nov  5 09:56 SUNWspro/
drwxrwxrwx  19 1000     1000          40 Nov  6 19:16 emacs-22.1/
```

Once a snapshot has been created, to synchronize the file system again, you create a new snapshot and then use the incremental snapshot feature of `zfs send` to send the changes between the two snapshots to the slave machine again:

```
root-shell> zfs send -i scratchpool@snapshot1 scratchpool@snapshot2 |ssh id@host pfexec zfs recv slavep
```

This operation only succeeds if the file system on the slave machine has not been modified at all. You cannot apply the incremental changes to a destination file system that has changed. In the example above, the `ls` command would cause problems by changing the metadata, such as the last access time for files or directories.

To prevent changes on the slave file system, set the file system on the slave to be read-only:

```
root-shell> zfs set readonly=on slavepool
```

Setting `readonly` means that you cannot change the file system on the slave by normal means, including the file system metadata. Operations that would normally update metadata (like our `ls`) silently perform their function without attempting to update the file system state.

In essence, the slave file system is nothing but a static copy of the original file system. However, even when configured to be read-only, a file system can have snapshots applied to it. With the file system set to read only, re-run the initial copy:

```
root-shell> zfs send scratchpool@snap1 |ssh id@host pfexec zfs recv -F slavepool
```

Now you can make changes to the original file system and replicate them to the slave.

17.1.2 Configuring MySQL for ZFS Replication

Configuring MySQL on the source file system is a case of creating the data on the file system that you intend to replicate. The configuration file in the example below has been updated to use `/scratchpool/mysql-data` as the data directory, and now you can initialize the tables:

```
root-shell> mysql_install_db --defaults-file=/etc/mysql/5.5/my.cnf --user=mysql
```

To synchronize the initial information, perform a new snapshot and then send an incremental snapshot to the slave using `zfs send`:

```
root-shell> zfs snapshot scratchpool@snap2
root-shell> zfs send -i scratchpool@snap1 scratchpool@snap2|ssh id@host pfexec zfs recv slavepool
```

Doublecheck that the slave has the data by looking at the MySQL data directory on the `slavepool`:

```
root-shell> ls -al /slavepool/mysql-data/
```

Now you can start up MySQL, create some data, and then replicate the changes using `zfs send`/`zfs recv` to the slave to synchronize the changes.

The rate at which you perform the synchronization depends on your application and environment. The limitation is the speed required to perform the snapshot and then to send the changes over the network.

To automate the process, create a script that performs the snapshot, send, and receive operation, and use `cron` to synchronize the changes at set times or intervals.

17.1.3 Handling MySQL Recovery with ZFS

When using ZFS replication to provide a constant copy of your data, ensure that you can recover your tables, either manually or automatically, in the event of a failure of the original system.

In the event of a failure, follow this sequence:

1. Stop the script on the master, if it is still up and running.
2. Set the slave file system to be read/write:

```
root-shell> zfs set readonly=off slavepool
```

3. Start up `mysqld` on the slave. If you are using `InnoDB`, you get auto-recovery, if it is needed, to make sure the table data is correct, as shown here when I started up from our mid-INSERT snapshot:

```
InnoDB: The log sequence number in ibdata files does not match
InnoDB: the log sequence number in the ib_logfiles!
081109 15:59:59  InnoDB: Database was not shut down normally!
InnoDB: Starting crash recovery.
InnoDB: Reading tablespace information from the .ibd files...
InnoDB: Restoring possible half-written data pages from the doublewrite
InnoDB: buffer...
081109 16:00:03  InnoDB: Started; log sequence number 0 1142807951
081109 16:00:03 [Note] /slavepool/mysql-5.0.67-solaris10-i386/bin/mysqld: ready for connections.
Version: '5.0.67'  socket: '/tmp/mysql.sock'  port: 3306  MySQL Community Server (GPL)
```

Use `InnoDB` tables and a regular synchronization schedule to reduce the risk for significant data loss. On MyISAM tables, you might need to run `REPAIR TABLE`, and you might even have lost some information.

17.2 Using MySQL with `memcached`

`memcached` is a simple, highly scalable key-based cache that stores data and objects wherever dedicated or spare RAM is available for quick access by applications, without going through layers of parsing or disk I/O. To use, you run the `memcached` command on one or more hosts and then use the shared cache to store objects. For more usage instructions, see Section 17.2.2, “Using `memcached`”

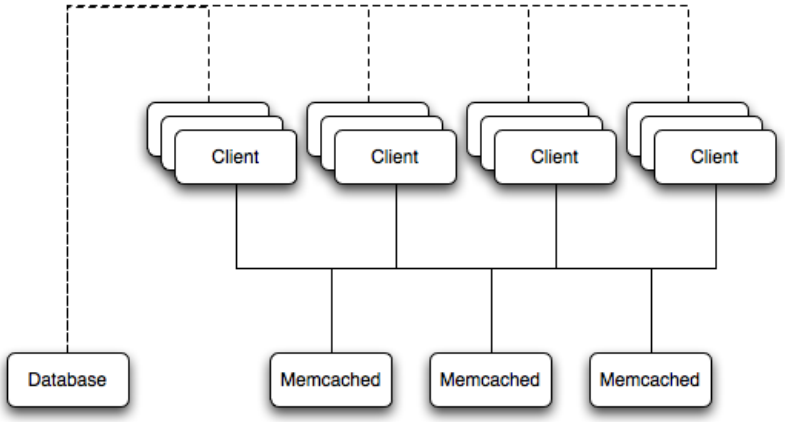
Benefits of using `memcached` include:

- Because all information is stored in RAM, the access speed is faster than loading the information each time from disk.
- Because the “value” portion of the key-value pair does not have any data type restrictions, you can cache data such as complex structures, documents, images, or a mixture of such things.
- If you use the in-memory cache to hold transient information, or as a read-only cache for information also stored in a database, the failure of any `memcached` server is not critical. For persistent data, you can fall back to an alternative lookup method using database queries, and reload the data into RAM on a different server.

The typical usage environment is to modify your application so that information is read from the cache provided by `memcached`. If the information is not in `memcached`, then the data is loaded from the MySQL database and written into the cache so that future requests for the same object benefit from the cached data.

For a typical deployment layout, see Figure 17.2, “`memcached` Architecture Overview”.

Figure 17.2 memcached Architecture Overview



In the example structure, any of the clients can contact one of the memcached servers to request a given key. Each client is configured to talk to all of the servers shown in the illustration. Within the client, when the request is made to store the information, the key used to reference the data is hashed and this hash is then used to select one of the memcached servers. The selection of the memcached server takes place on the client before the server is contacted, keeping the process lightweight.

The same algorithm is used again when a client requests the same key. The same key generates the same hash, and the same memcached server is selected as the source for the data. Using this method, the cached data is spread among all of the memcached servers, and the cached information is accessible from any client. The result is a distributed, memory-based, cache that can return information, particularly complex data and structures, much faster than natively reading the information from the database.

The data held within a traditional memcached server is never stored on disk (only in RAM, which means there is no persistence of data), and the RAM cache is always populated from the backing store (a MySQL database). If a memcached server fails, the data can always be recovered from the MySQL database.

17.2.1 Installing memcached

You can build and install memcached from the source code directly, or you can use an existing operating system package or installation.

Installing memcached from a Binary Distribution

To install memcached on a Red Hat, or Fedora host, use yum:

```
root-shell> yum install memcached
```

Note

On CentOS, you may be able to obtain a suitable RPM from another source, or use the source tarball.

To install memcached on a Debian or Ubuntu host, use apt-get:

```
root-shell> apt-get install memcached
```

To install memcached on a Gentoo host, use emerge:

```
root-shell> emerge install memcached
```

Building memcached from Source

On other Unix-based platforms, including Solaris, AIX, HP-UX and OS X, and Linux distributions not mentioned already, you must install from source. For Linux, make sure you have a 2.6-based kernel, which includes the improved epoll interface. For all platforms, ensure that you have libevent 1.1 or higher installed. You can obtain libevent from libevent web page.

You can obtain the source for memcached from memcached Web site.

To build memcached, follow these steps:

1. Extract the memcached source package:

```
shell> gunzip -c memcached-1.2.5.tar.gz | tar xf -
```

2. Change to the memcached-1.2.5 directory:

```
shell> cd memcached-1.2.5
```

3. Run configure

```
shell> ./configure
```

Some additional options you might specify to the configure:

- --prefix

  To specify a different installation directory, use the --prefix option:

  ```
  shell> ./configure --prefix=/opt
  ```

  The default is to use the /usr/local directory.

- --with-libevent

  If you have installed libevent and configure cannot find the library, use the --with-libevent option to specify the location of the installed library.

- --enable-64bit

  To build a 64-bit version of memcached (which enables you to use a single instance with a large RAM allocation), use --enable-64bit.

- --enable-threads

  To enable multi-threading support in memcached, which improves the response times on servers with a heavy load, use --enable-threads. You must have support for the POSIX threads within your operating system to enable thread support. For more information on the threading support, see Section 17.2.2.8, “memcached Thread Support”.

- --enable-dtrace

  memcached includes a range of DTrace threads that can be used to monitor and benchmark a memcached instance. For more information, see Section 17.2.2.6, “Using memcached and DTrace”.

4. Run make to build memcached:

```
shell> make
```

5. Run `make install` to install `memcached`:

```
shell> make install
```

17.2.2 Using memcached

To start using `memcached`, start the `memcached` service on one or more servers. Running `memcached` sets up the server, allocates the memory and starts listening for connections from clients.

Note

You do not need to be a privileged user (`root`) to run `memcached` except to listen on one of the privileged TCP/IP ports (below 1024). You must, however, use a user that has not had their memory limits restricted using `setrlimit` or similar.

To start the server, run `memcached` as a nonprivileged (that is, non-`root`) user:

```
shell> memcached
```

By default, `memcached` uses the following settings:

- Memory allocation of 64MB
- Listens for connections on all network interfaces, using port 11211
- Supports a maximum of 1024 simultaneous connections

Typically, you would specify the full combination of options that you want when starting `memcached`, and normally provide a startup script to handle the initialization of `memcached`. For example, the following line starts `memcached` with a maximum of 1024MB RAM for the cache, listening on port 11211 on the IP address 192.168.0.110, running as a background daemon:

```
shell> memcached -d -m 1024 -p 11211 -l 192.168.0.110
```

To ensure that `memcached` is started up on boot, check the init script and configuration parameters.

17.2.2.1 memcached Command-Line Options

`memcached` supports the following options:

- `-u user`

  If you start `memcached` as `root`, use the `-u` option to specify the user for executing `memcached`:

  ```
  shell> memcached -u memcache
  ```

- `-m memory`

  Set the amount of memory allocated to `memcached` for object storage. Default is 64MB.

  To increase the amount of memory allocated for the cache, use the `-m` option to specify the amount of RAM to be allocated (in megabytes). The more RAM you allocate, the more data you can store and therefore the more effective your cache is.

Warning

Do not specify a memory allocation larger than your available RAM. If you specify too large a value, then some RAM allocated for `memcached` uses swap space, and not physical RAM. This may lead to delays when storing

and retrieving values, because data is swapped to disk, instead of storing the data directly in RAM.

You can use the output of the vmstat command to get the free memory, as shown in free column:

```
shell> vmstat
kthr      memory            page            disk          faults      cpu
r b w   swap  free  re  mf pi po fr de sr s1 s2 -- --   in   sy   cs us sy id
0 0 0 5170504 3450392 2  7  2  0  0  0  4  0  0  0  0  296   54  199  0  0 100
```

For example, to allocate 3GB of RAM:

```
shell> memcached -m 3072
```

On 32-bit x86 systems where you are using PAE to access memory above the 4GB limit, you cannot allocate RAM beyond the maximum process size. You can get around this by running multiple instances of memcached, each listening on a different port:

```
shell> memcached -m 1024 -p11211
shell> memcached -m 1024 -p11212
shell> memcached -m 1024 -p11213
```

Note

On all systems, particularly 32-bit, ensure that you leave enough room for both memcached application in addition to the memory setting. For example, if you have a dedicated memcached host with 4GB of RAM, do not set the memory size above 3500MB. Failure to do this may cause either a crash or severe performance issues.

- `-l interface`

  Specify a network interface/address to listen for connections. The default is to listen on all available address (INADDR_ANY).

  ```
  shell> memcached -l 192.168.0.110
  ```

  Support for IPv6 address support was added in memcached 1.2.5.

- `-p port`

  Specify the TCP port to use for connections. Default is 18080.

  ```
  shell> memcached -p 18080
  ```

- `-U port`

  Specify the UDP port to use for connections. Default is 11211, 0 switches UDP off.

  ```
  shell> memcached -U 18080
  ```

- `-s socket`

  Specify a Unix socket to listen on.

  If you are running memcached on the same server as the clients, you can disable the network interface and use a local Unix socket using the -s option:

  ```
  ```

```
shell> memcached -s /tmp/memcached
```

Using a Unix socket automatically disables network support, and saves network ports (allowing more ports to be used by your web server or other process).

- `-a mask`

  Specify the access mask to be used for the Unix socket, in octal. Default is 0700.

- `-c connections`

  Specify the maximum number of simultaneous connections to the `memcached` service. The default is 1024.

  ```
  shell> memcached -c 2048
  ```

  Use this option, either to reduce the number of connections (to prevent overloading `memcached` service) or to increase the number to make more effective use of the server running `memcached` server.

- `-t threads`

  Specify the number of threads to use when processing incoming requests.

  By default, `memcached` is configured to use 4 concurrent threads. The threading improves the performance of storing and retrieving data in the cache, using a locking system to prevent different threads overwriting or updating the same values. To increase or decrease the number of threads, use the `-t` option:

  ```
  shell> memcached -t 8
  ```

- `-d`

  Run `memcached` as a daemon (background) process:

  ```
  shell> memcached -d
  ```

- `-r`

  Maximize the size of the core file limit. In the event of a failure, this attempts to dump the entire memory space to disk as a core file, up to any limits imposed by `setrlimit`.

- `-M`

  Return an error to the client when the memory has been exhausted. This replaces the normal behavior of removing older items from the cache to make way for new items.

- `-k`

  Lock down all paged memory. This reserves the memory before use, instead of allocating new slabs of memory as new items are stored in the cache.

  Note

  There is a user-level limit on how much memory you can lock. Trying to allocate more than the available memory fails. You can set the limit for the user you started the daemon with (not for the `-u user` user) within the shell by using `ulimit -S -l NUM_KB`

- `-v`

Verbose mode. Prints errors and warnings while executing the main event loop.

- `-vv`

  Very verbose mode. In addition to information printed by `-v`, also prints each client command and the response.

- `-vvv`

  Extremely verbose mode. In addition to information printed by `-vv`, also show the internal state transitions.

- `-h`

  Print the help message and exit.

- `-i`

  Print the `memcached` and `libevent` license.

- `-I mem`

  Specify the maximum size permitted for storing an object within the `memcached` instance. The size supports a unit postfix (`k` for kilobytes, `m` for megabytes). For example, to increase the maximum supported object size to 32MB:

  ```
  shell> memcached -I 32m
  ```

  The maximum object size you can specify is 128MB, the default remains at 1MB.

  This option was added in 1.4.2.

- `-b`

  Set the backlog queue limit. The backlog queue configures how many network connections can be waiting to be processed by `memcached`. Increasing this limit may reduce errors received by the client that it is not able to connect to the `memcached` instance, but does not improve the performance of the server. The default is 1024.

- `-P pidfile`

  Save the process ID of the `memcached` instance into `file`.

- `-f`

  Set the chunk size growth factor. When allocating new memory chunks, the allocated size of new chunks is determined by multiplying the default slab size by this factor.

  To see the effects of this option without extensive testing, use the `-vv` command-line option to show the calculated slab sizes. For more information, see Section 17.2.2.9, “`memcached` Logs”.

- `-n bytes`

  The minimum space allocated for the key+value+flags information. The default is 48 bytes.

- `-L`

  On systems that support large memory pages, enables large memory page use. Using large memory pages enables `memcached` to allocate the item cache in one large chunk, which can improve the performance by reducing the number misses when accessing memory.

- `-C`

Disable the use of compare and swap (CAS) operations.

This option was added in memcached 1.3.x.

- -D char

  Set the default character to be used as a delimiter between the key prefixes and IDs. This is used for the per-prefix statistics reporting (see Section 17.2.4, “Getting memcached Statistics”). The default is the colon (:). If this option is used, statistics collection is turned on automatically. If not used, you can enable stats collection by sending the stats detail on command to the server.

  This option was added in memcached 1.3.x.

- -R num

  Sets the maximum number of requests per event process. The default is 20.

- -B protocol

  Set the binding protocol, that is, the default memcached protocol support for client connections. Options are ascii, binary or auto. Automatic (auto) is the default.

  This option was added in memcached 1.4.0.

17.2.2.2 memcached Deployment

When using memcached you can use a number of different potential deployment strategies and topologies. The exact strategy to use depends on your application and environment. When developing a system for deploying memcached within your system, keep in mind the following points:

- memcached is only a caching mechanism. It shouldn't be used to store information that you cannot otherwise afford to lose and then load from a different location.
- There is no security built into the memcached protocol. At a minimum, make sure that the servers running memcached are only accessible from inside your network, and that the network ports being used are blocked (using a firewall or similar). If the information on the memcached servers that is being stored is any sensitive, then encrypt the information before storing it in memcached.
- memcached does not provide any sort of failover. Because there is no communication between different memcached instances. If an instance fails, your application must capable of removing it from the list, reloading the data and then writing data to another memcached instance.
- Latency between the clients and the memcached can be a problem if you are using different physical machines for these tasks. If you find that the latency is a problem, move the memcached instances to be on the clients.
- Key length is determined by the memcached server. The default maximum key size is 250 bytes.
- Try to use at least two memcached instances, especially for multiple clients, to avoid having a single point of failure. Ideally, create as many memcached nodes as possible. When adding and removing memcached instances from a pool, the hashing and distribution of key/value pairs may be affected. For information on how to avoid problems, see Section 17.2.2.5, “memcached Hashing/Distribution Types”.

17.2.2.3 Using Namespaces

The memcached cache is a very simple massive key/value storage system, and as such there is no way of compartmentalizing data automatically into different sections. For example, if you are storing information by the unique ID returned from a MySQL database, then storing the data from two different tables could run into issues because the same ID might be valid in both tables.

Some interfaces provide an automated mechanism for creating *namespaces* when storing information into the cache. In practice, these namespaces are merely a prefix before a given ID that is applied every time a value is stored or retrieve from the cache.

You can implement the same basic principle by using keys that describe the object and the unique identifier within the key that you supply when the object is stored. For example, when storing user data, prefix the ID of the user with `user:` or `user-`.

Note

Using namespaces or prefixes only controls the keys stored/retrieved. There is no security within `memcached`, and therefore no way to enforce that a particular client only accesses keys with a particular namespace. Namespaces are only useful as a method of identifying data and preventing corruption of key/value pairs.

17.2.2.4 Data Expiry

There are two types of data expiry within a `memcached` instance. The first type is applied at the point when you store a new key/value pair into the `memcached` instance. If there is not enough space within a suitable slab to store the value, then an existing least recently used (LRU) object is removed (evicted) from the cache to make room for the new item.

The LRU algorithm ensures that the object that is removed is one that is either no longer in active use or that was used so long ago that its data is potentially out of date or of little value. However, in a system where the memory allocated to `memcached` is smaller than the number of regularly used objects required in the cache, a lot of expired items could be removed from the cache even though they are in active use. You use the statistics mechanism to get a better idea of the level of evictions (expired objects). For more information, see Section 17.2.4, “Getting `memcached` Statistics”.

You can change this eviction behavior by setting the `-M` command-line option when starting `memcached`. This option forces an error to be returned when the memory has been exhausted, instead of automatically evicting older data.

The second type of expiry system is an explicit mechanism that you can set when a key/value pair is inserted into the cache, or when deleting an item from the cache. Using an expiration time can be a useful way of ensuring that the data in the cache is up to date and in line with your application needs and requirements.

A typical scenario for explicitly setting the expiry time might include caching session data for a user when accessing a Web site. `memcached` uses a lazy expiry mechanism where the explicit expiry time that has been set is compared with the current time when the object is requested. Only objects that have not expired are returned.

You can also set the expiry time when explicitly deleting an object from the cache. In this case, the expiry time is really a timeout and indicates the period when any attempts to set the value for a given key are rejected.

17.2.2.5 `memcached` Hashing/Distribution Types

The `memcached` client interface supports a number of different distribution algorithms that are used in multi-server configurations to determine which host should be used when setting or getting data from a given `memcached` instance. When you get or set a value, a hash is constructed from the supplied key and then used to select a host from the list of configured servers. Because the hashing mechanism uses the supplied key as the basis for the hash, the same server is selected during both set and get operations.

You can think of this process as follows. Given an array of servers (a, b, and c), the client uses a hashing algorithm that returns an integer based on the key being stored or retrieved. The resulting value is then used to select a server from the list of servers configured in the client. Most standard

client hashing within memcache clients uses a simple modulus calculation on the value against the number of configured memcached servers. You can summarize the process in pseudocode as:

```
@memcservers = ['a.memc','b.memc','c.memc'];
$value = hash($key);
$chosen = $value % length(@memcservers);
```

Replacing the above with values:

```
@memcservers = ['a.memc','b.memc','c.memc'];
$value = hash('myid');
$chosen = 7009 % 3;
```

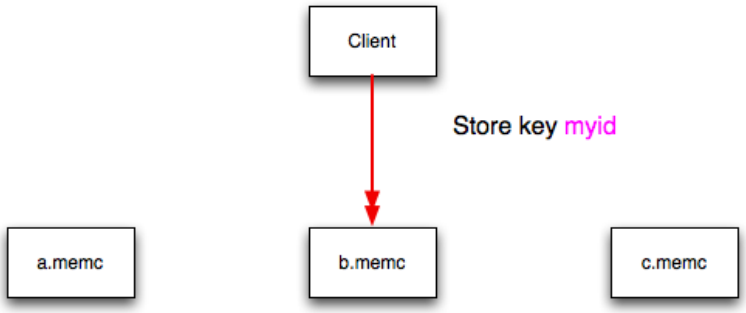
In the above example, the client hashing algorithm chooses the server at index 1 (`7009 % 3 = 1`), and store or retrieve the key and value with that server.

Note

This selection and hashing process is handled automatically by the memcached client you are using; you need only provide the list of memcached servers to use.

You can see a graphical representation of this below in Figure 17.3, "memcached Hash Selection".

Figure 17.3 memcached Hash Selection



The same hashing and selection process takes place during any operation on the specified key within the memcached client.

Using this method provides a number of advantages:

- The hashing and selection of the server to contact is handled entirely within the client. This eliminates the need to perform network communication to determine the right machine to contact.
- Because the determination of the memcached server occurs entirely within the client, the server can be selected automatically regardless of the operation being executed (set, get, increment, etc.).
- Because the determination is handled within the client, the hashing algorithm returns the same value for a given key; values are not affected or reset by differences in the server environment.
- Selection is very fast. The hashing algorithm on the key value is quick and the resulting selection of the server is from a simple array of available machines.
- Using client-side hashing simplifies the distribution of data over each memcached server. Natural distribution of the values returned by the hashing algorithm means that keys are automatically spread over the available servers.

Providing that the list of servers configured within the client remains the same, the same stored key returns the same value, and therefore selects the same server.

However, if you do not use the same hashing mechanism then the same data may be recorded on different servers by different interfaces, both wasting space on your memcached and leading to potential differences in the information.

Note

One way to use a multi-interface compatible hashing mechanism is to use the libmemcached library and the associated interfaces. Because the interfaces for the different languages (including C, Ruby, Perl and Python) use the same client library interface, they always generate the same hash code from the ID.

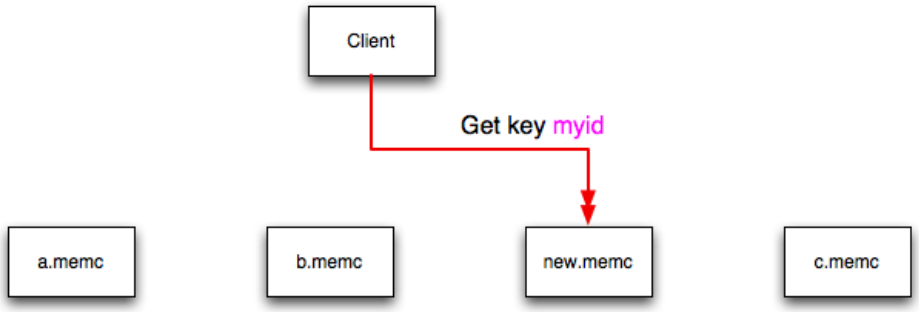
The problem with client-side selection of the server is that the list of the servers (including their sequential order) *must* remain consistent on each client using the memcached servers, and the servers must be available. If you try to perform an operation on a key when:

- A new memcached instance has been added to the list of available instances
- A memcached instance has been removed from the list of available instances
- The order of the memcached instances has changed

When the hashing algorithm is used on the given key, but with a different list of servers, the hash calculation may choose a different server from the list.

If a new memcached instance is added into the list of servers, as new.memc is in the example below, then a GET operation using the same key, myid, can result in a cache-miss. This is because the same value is computed from the key, which selects the same index from the array of servers, but index 2 now points to the new server, not the server c.memc where the data was originally stored. This would result in a cache miss, even though the key exists within the cache on another memcached instance.

Figure 17.4 memcached Hash Selection with New memcached instance



This means that servers c.memc and new.memc both contain the information for key myid, but the information stored against the key in eachs server may be different in each instance. A more significant problem is a much higher number of cache-misses when retrieving data, as the addition of a new server changes the distribution of keys, and this in turn requires rebuilding the cached data on the memcached instances, causing an increase in database reads.

The same effect can occur if you actively manage the list of servers configured in your clients, adding and removing the configured memcached instances as each instance is identified as being available. For example, removing a memcached instance when the client notices that the instance can no longer be contacted can cause the server selection to fail as described here.

To prevent this causing significant problems and invalidating your cache, you can select the hashing algorithm used to select the server. There are two common types of hashing algorithm, *consistent* and *modula*.

With *consistent* hashing algorithms, the same key when applied to a list of servers always uses the same server to store or retrieve the keys, even if the list of configured servers changes. This means

that you can add and remove servers from the configure list and always use the same server for a given key. There are two types of consistent hashing algorithms available, Ketama and Wheel. Both types are supported by libmemcached, and implementations are available for PHP and Java.

Any consistent hashing algorithm has some limitations. When you add servers to an existing list of configured servers, keys are distributed to the new servers as part of the normal distribution. When you remove servers from the list, the keys are re-allocated to another server within the list, meaning that the cache needs to be re-populated with the information. Also, a consistent hashing algorithm does not resolve the issue where you want consistent selection of a server across multiple clients, but where each client contains a different list of servers. The consistency is enforced only within a single client.

With a *modula* hashing algorithm, the client selects a server by first computing the hash and then choosing a server from the list of configured servers. As the list of servers changes, so the server selected when using a modula hashing algorithm also changes. The result is the behavior described above; changes to the list of servers mean that different servers are selected when retrieving data, leading to cache misses and increase in database load as the cache is re-seeded with information.

If you use only a single memcached instance for each client, or your list of memcached servers configured for a client never changes, then the selection of a hashing algorithm is irrelevant, as it has no noticeable effect.

If you change your servers regularly, or you use a common set of servers that are shared among a large number of clients, then using a consistent hashing algorithm should help to ensure that your cache data is not duplicated and the data is evenly distributed.

17.2.2.6 Using memcached and DTrace

memcached includes a number of different DTrace probes that can be used to monitor the operation of the server. The probes included can monitor individual connections, slab allocations, and modifications to the hash table when a key/value pair is added, updated, or removed.

For more information on DTrace and writing DTrace scripts, read the DTrace User Guide.

Support for DTrace probes was added to memcached 1.2.6 includes a number of DTrace probes that can be used to help monitor your application. DTrace is supported on Solaris 10, OpenSolaris, OS X 10.5 and FreeBSD. To enable the DTrace probes in memcached, build from source and use the --enable-dtrace option. For more information, see Section 17.2.1, “Installing memcached”.

The probes supported by memcached are:

- conn-allocate(connid)

  Fired when a connection object is allocated from the connection pool.

  - connid: The connection ID.

- conn-release(connid)

  Fired when a connection object is released back to the connection pool.

  Arguments:

  - connid: The connection ID.

- conn-create(ptr)

  Fired when a new connection object is being created (that is, there are no free connection objects in the connection pool).

  Arguments:

  - ptr: A pointer to the connection. object

- `conn-destroy(ptr)`

  Fired when a connection object is being destroyed.

  Arguments:

  - `ptr`: A pointer to the connection object.

- `conn-dispatch(connid, threadid)`

  Fired when a connection is dispatched from the main or connection-management thread to a worker thread.

  Arguments:

  - `connid`: The connection ID.
  - `threadid`: The thread ID.

- `slabs-allocate(size, slabclass, slabsize, ptr)`

  Allocate memory from the slab allocator.

  Arguments:

  - `size`: The requested size.
  - `slabclass`: The allocation is fulfilled in this class.
  - `slabsize`: The size of each item in this class.
  - `ptr`: A pointer to allocated memory.

- `slabs-allocate-failed(size, slabclass)`

  Failed to allocate memory (out of memory).

  Arguments:

  - `size`: The requested size.
  - `slabclass`: The class that failed to fulfill the request.

- `slabs-slabclass-allocate(slabclass)`

  Fired when a slab class needs more space.

  Arguments:

  - `slabclass`: The class that needs more memory.

- `slabs-slabclass-allocate-failed(slabclass)`

  Failed to allocate memory (out of memory).

  Arguments:

  - `slabclass`: The class that failed to grab more memory.

- `slabs-free(size, slabclass, ptr)`

  Release memory.

  Arguments:

  - size: The amount of memory to release, in bytes.
  - slabclass: The class the memory belongs to.
  - ptr: A pointer to the memory to release.
- assoc-find(key, depth)

  Fired when we have searched the hash table for a named key. These two elements provide an insight into how well the hash function operates. Traversals are a sign of a less optimal function, wasting CPU capacity.

  Arguments:

  - key: The key searched for.
  - depth: The depth in the list of hash table.
- assoc-insert(key, nokeys)

  Fired when a new item has been inserted.

  Arguments:

  - key: The key just inserted.
  - nokeys: The total number of keys currently being stored, including the key for which insert was called.
- assoc-delete(key, nokeys)

  Fired when a new item has been removed.

  Arguments:

  - key: The key just deleted.
  - nokeys: The total number of keys currently being stored, excluding the key for which delete was called.
- item-link(key, size)

  Fired when an item is being linked in the cache.

  Arguments:

  - key: The items key.
  - size: The size of the data.
- item-unlink(key, size)

  Fired when an item is being deleted.

  Arguments:

  - key: The items key.
  - size: The size of the data.
- item-remove(key, size)

  Fired when the refcount for an item is reduced.

Arguments:

- key: The item's key.
- size: The size of the data.

- item-update(key, size)

  Fired when the "last referenced" time is updated.

  Arguments:

  - key: The item's key.
  - size: The size of the data.

- item-replace(oldkey, oldsize, newkey, newsize)

  Fired when an item is being replaced with another item.

  Arguments:

  - oldkey: The key of the item to replace.
  - oldsize: The size of the old item.
  - newkey: The key of the new item.
  - newsize: The size of the new item.

- process-command-start(connid, request, size)

  Fired when the processing of a command starts.

  Arguments:

  - connid: The connection ID.
  - request: The incoming request.
  - size: The size of the request.

- process-command-end(connid, response, size)

  Fired when the processing of a command is done.

  Arguments:

  - connid: The connection ID.
  - response: The response to send back to the client.
  - size: The size of the response.

- command-get(connid, key, size)

  Fired for a get command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.

    - size: The size of the key's data (or -1 if not found).
- command-gets(connid, key, size, casid)

  Fired for a gets command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.
  - size: The size of the key's data (or -1 if not found).
  - casid: The casid for the item.
- command-add(connid, key, size)

  Fired for a add command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.
  - size: The new size of the key's data (or -1 if not found).
- command-set(connid, key, size)

  Fired for a set command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.
  - size: The new size of the key's data (or -1 if not found).
- command-replace(connid, key, size)

  Fired for a replace command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.
  - size: The new size of the key's data (or -1 if not found).
- command-prepend(connid, key, size)

  Fired for a prepend command.

  Arguments:

  - connid: The connection ID.
  - key: The requested key.
  - size: The new size of the key's data (or -1 if not found).

- `command-append(connid, key, size)`

  Fired for a `append` command.

  Arguments:

  - `connid`: The connection ID.
  - `key`: The requested key.
  - `size`: The new size of the key's data (or -1 if not found).

- `command-cas(connid, key, size, casid)`

  Fired for a `cas` command.

  Arguments:

  - `connid`: The connection ID.
  - `key`: The requested key.
  - `size`: The size of the key's data (or -1 if not found).
  - `casid`: The cas ID requested.

- `command-incr(connid, key, val)`

  Fired for `incr` command.

  Arguments:

  - `connid`: The connection ID.
  - `key`: The requested key.
  - `val`: The new value.

- `command-decr(connid, key, val)`

  Fired for `decr` command.

  Arguments:

  - `connid`: The connection ID.
  - `key`: The requested key.
  - `val`: The new value.

- `command-delete(connid, key, exptime)`

  Fired for a `delete` command.

  Arguments:

  - `connid`: The connection ID.
  - `key`: The requested key.
  - `exptime`: The expiry time.

17.2.2.7 Memory Allocation within memcached

When you first start memcached, the memory that you have configured is not automatically allocated. Instead, memcached only starts allocating and reserving physical memory once you start saving information into the cache.

When you start to store data into the cache, memcached does not allocate the memory for the data on an item by item basis. Instead, a slab allocation is used to optimize memory usage and prevent memory fragmentation when information expires from the cache.

With slab allocation, memory is reserved in blocks of 1MB. The slab is divided up into a number of blocks of equal size. When you try to store a value into the cache, memcached checks the size of the value that you are adding to the cache and determines which slab contains the right size allocation for the item. If a slab with the item size already exists, the item is written to the block within the slab.

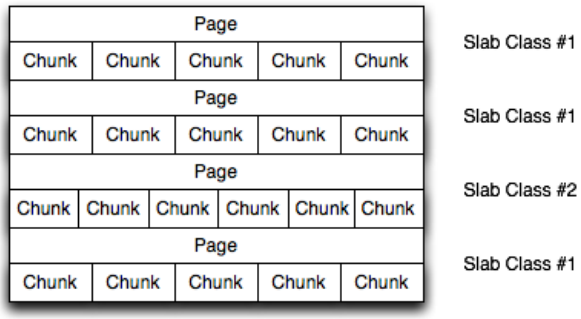
If the new item is bigger than the size of any existing blocks, then a new slab is created, divided up into blocks of a suitable size. If an existing slab with the right block size already exists, but there are no free blocks, a new slab is created. If you update an existing item with data that is larger than the existing block allocation for that key, then the key is re-allocated into a suitable slab.

For example, the default size for the smallest block is 88 bytes (40 bytes of value, and the default 48 bytes for the key and flag data). If the size of the first item you store into the cache is less than 40 bytes, then a slab with a block size of 88 bytes is created and the value stored.

If the size of the data that you intend to store is larger than this value, then the block size is increased by the chunk size factor until a block size large enough to hold the value is determined. The block size is always a function of the scale factor, rounded up to a block size which is exactly divisible into the chunk size.

For a sample of the structure, see Figure 17.5, “Memory Allocation in memcached”.

Figure 17.5 Memory Allocation in memcached



The result is that you have multiple pages allocated within the range of memory allocated to memcached. Each page is 1MB in size (by default), and is split into a different number of chunks, according to the chunk size required to store the key/value pairs. Each instance has multiple pages allocated, and a page is always created when a new item needs to be created requiring a chunk of a particular size. A slab may consist of multiple pages, and each page within a slab contains an equal number of chunks.

The chunk size of a new slab is determined by the base chunk size combined with the chunk size growth factor. For example, if the initial chunks are 104 bytes in size, and the default chunk size growth factor is used (1.25), then the next chunk size allocated would be the best power of 2 fit for 104*1.25, or 136 bytes.

Allocating the pages in this way ensures that memory does not get fragmented. However, depending on the distribution of the objects that you store, it may lead to an inefficient distribution of the slabs and chunks if you have significantly different sized items. For example, having a relatively small number of items within each chunk size may waste a lot of memory with just few chunks in each allocated page.

You can tune the growth factor to reduce this effect by using the `-f` command line option, which adapts the growth factor applied to make more effective use of the chunks and slabs allocated. For information on how to determine the current slab allocation statistics, see Section 17.2.4.2, "`memcached` Slabs Statistics".

If your operating system supports it, you can also start `memcached` with the `-L` command line option. This option preallocates all the memory during startup using large memory pages. This can improve performance by reducing the number of misses in the CPU memory cache.

17.2.2.8 `memcached` Thread Support

If you enable the thread implementation within when building `memcached` from source, then `memcached` uses multiple threads in addition to the `libevent` system to handle requests.

When enabled, the threading implementation operates as follows:

- Threading is handled by wrapping functions within the code to provide basic protection from updating the same global structures at the same time.
- Each thread uses its own instance of the `libevent` to help improve performance.
- TCP/IP connections are handled with a single thread listening on the TCP/IP socket. Each connection is then distributed to one of the active threads on a simple round-robin basis. Each connection then operates solely within this thread while the connection remains open.
- For UDP connections, all the threads listen to a single UDP socket for incoming requests. Threads that are not currently dealing with another request ignore the incoming packet. One of the remaining, nonbusy, threads reads the request and sends the response. This implementation can lead to increased CPU load as threads wake from sleep to potentially process the request.

Using threads can increase the performance on servers that have multiple CPU cores available, as the requests to update the hash table can be spread between the individual threads. To minimize overhead from the locking mechanism employed, experiment with different thread values to achieve the best performance based on the number and type of requests within your given workload.

17.2.2.9 `memcached` Logs

If you enable verbose mode, using the `-v`, `-vv`, or `-vvv` options, then the information output by `memcached` includes details of the operations being performed.

Without the verbose options, `memcached` normally produces no output during normal operating.

- **Output when using `-v`**

  The lowest verbosity level shows you:

  - Errors and warnings
  - Transient errors
  - Protocol and socket errors, including exhausting available connections
  - Each registered client connection, including the socket descriptor number and the protocol used.

    For example:

    ```
    32: Client using the ascii protocol
    33: Client using the ascii protocol
    ```

    The socket descriptor is only valid while the client remains connected. Non-persistent connections may not be effectively represented.

  Examples of the error messages output at this level include:

```
<%d send buffer was %d, now %d
Can't listen for events on fd %d
Can't read from libevent pipe
Catastrophic: event fd doesn't match conn fd!
Couldn't build response
Couldn't realloc input buffer
Couldn't update event
Failed to build UDP headers
Failed to read, and not due to blocking
Too many open connections
Unexpected state %d
```

- **Output when using `-vv`**

  When using the second level of verbosity, you get more detailed information about protocol operations, keys updated, chunk and network operatings and details.

  During the initial start-up of `memcached` with this level of verbosity, you are shown the sizes of the individual slab classes, the chunk sizes, and the number of entries per slab. These do not show the allocation of the slabs, just the slabs that would be created when data is added. You are also given information about the listen queues and buffers used to send information. A sample of the output generated for a TCP/IP based system with the default memory and growth factors is given below:

```
shell> memcached -vv
slab class   1: chunk size     80 perslab 13107
slab class   2: chunk size    104 perslab 10082
slab class   3: chunk size    136 perslab  7710
slab class   4: chunk size    176 perslab  5957
slab class   5: chunk size    224 perslab  4681
slab class   6: chunk size    280 perslab  3744
slab class   7: chunk size    352 perslab  2978
slab class   8: chunk size    440 perslab  2383
slab class   9: chunk size    552 perslab  1899
slab class  10: chunk size    696 perslab  1506
slab class  11: chunk size    872 perslab  1202
slab class  12: chunk size   1096 perslab   956
slab class  13: chunk size   1376 perslab   762
slab class  14: chunk size   1720 perslab   609
slab class  15: chunk size   2152 perslab   487
slab class  16: chunk size   2696 perslab   388
slab class  17: chunk size   3376 perslab   310
slab class  18: chunk size   4224 perslab   248
slab class  19: chunk size   5280 perslab   198
slab class  20: chunk size   6600 perslab   158
slab class  21: chunk size   8256 perslab   127
slab class  22: chunk size  10320 perslab   101
slab class  23: chunk size  12904 perslab    81
slab class  24: chunk size  16136 perslab    64
slab class  25: chunk size  20176 perslab    51
slab class  26: chunk size  25224 perslab    41
slab class  27: chunk size  31536 perslab    33
slab class  28: chunk size  39424 perslab    26
slab class  29: chunk size  49280 perslab    21
slab class  30: chunk size  61600 perslab    17
slab class  31: chunk size  77000 perslab    13
slab class  32: chunk size  96256 perslab    10
slab class  33: chunk size 120320 perslab     8
slab class  34: chunk size 150400 perslab     6
slab class  35: chunk size 188000 perslab     5
slab class  36: chunk size 235000 perslab     4
slab class  37: chunk size 293752 perslab     3
slab class  38: chunk size 367192 perslab     2
slab class  39: chunk size 458992 perslab     2
<26 server listening (auto-negotiate)
<29 server listening (auto-negotiate)
<30 send buffer was 57344, now 2097152
<31 send buffer was 57344, now 2097152
```

```
<30 server listening (udp)
<30 server listening (udp)
<31 server listening (udp)
<30 server listening (udp)
<30 server listening (udp)
<31 server listening (udp)
<31 server listening (udp)
<31 server listening (udp)
```

Using this verbosity level can be a useful way to check the effects of the growth factor used on slabs with different memory allocations, which in turn can be used to better tune the growth factor to suit the data you are storing in the cache. For example, if you set the growth factor to 4 (quadrupling the size of each slab):

```
shell> memcached -f 4 -m 1g -vv
slab class   1: chunk size     80 perslab 13107
slab class   2: chunk size    320 perslab  3276
slab class   3: chunk size   1280 perslab   819
slab class   4: chunk size   5120 perslab   204
slab class   5: chunk size  20480 perslab    51
slab class   6: chunk size  81920 perslab    12
slab class   7: chunk size 327680 perslab     3
...
```

During use of the cache, this verbosity level also prints out detailed information on the storage and recovery of keys and other information. An example of the output during a typical set/get and increment/decrement operation is shown below.

```
32: Client using the ascii protocol
<32 set my_key 0 0 10
>32 STORED
<32 set object_key 1 0 36
>32 STORED
<32 get my_key
>32 sending key my_key
>32 END
<32 get object_key
>32 sending key object_key
>32 END
<32 set key 0 0 6
>32 STORED
<32 incr key 1
>32 789544
<32 decr key 1
>32 789543
<32 incr key 2
>32 789545
<32 set my_key 0 0 10
>32 STORED
<32 set object_key 1 0 36
>32 STORED
<32 get my_key
>32 sending key my_key
>32 END
<32 get object_key
>32 sending key object_key1 1 36

>32 END
<32 set key 0 0 6
>32 STORED
<32 incr key 1
>32 789544
<32 decr key 1
>32 789543
<32 incr key 2
>32 789545
```

During client communication, for each line, the initial character shows the direction of flow of the information. The < for communication from the client to the memcached server and > for communication back to the client. The number is the numeric socket descriptor for the connection.

- **Output when using -vvv**

  This level of verbosity includes the transitions of connections between different states in the event library while reading and writing content to/from the clients. It should be used to diagnose and identify issues in client communication. For example, you can use this information to determine if memcached is taking a long time to return information to the client, during the read of the client operation or before returning and completing the operation. An example of the typical sequence for a set operation is provided below:

```
<32 new auto-negotiating client connection
32: going from conn_new_cmd to conn_waiting
32: going from conn_waiting to conn_read
32: going from conn_read to conn_parse_cmd
32: Client using the ascii protocol
<32 set my_key 0 0 10
32: going from conn_parse_cmd to conn_nread
> NOT FOUND my_key
>32 STORED
32: going from conn_nread to conn_write
32: going from conn_write to conn_new_cmd
32: going from conn_new_cmd to conn_waiting
32: going from conn_waiting to conn_read
32: going from conn_read to conn_closing
<32 connection closed.
```

All of the verbosity levels in memcached are designed to be used during debugging or examination of issues. The quantity of information generated, particularly when using -vvv, is significant, particularly on a busy server. Also be aware that writing the error information out, especially to disk, may negate some of the performance gains you achieve by using memcached. Therefore, use in production or deployment environments is not recommended.

17.2.3 Developing a memcached Application

A number of language interfaces let applications store and retrieve information with memcached servers. You can write memcached applications in popular languages such as Perl, PHP, Python, Ruby, C, and Java.

Data stored into a memcached server is referred to by a single string (the key), with storage into the cache and retrieval from the cache using the key as the reference. The cache therefore operates like a large associative array or hash table. It is not possible to structure or otherwise organize the information stored in the cache. To emulate database notions such as multiple tables or composite key values, you must encode the extra information into the strings used as keys. For example, to store or look up the address corresponding to a specific latitude and longitude, you might turn those two numeric values into a single comma-separated string to use as a key.

17.2.3.1 Basic memcached Operations

The interface to memcached supports the following methods for storing and retrieving information in the cache, and these are consistent across all the different APIs, although the language specific mechanics might be different:

- get(*key*): Retrieves information from the cache. Returns the value associated with the key if the specified key exists. Returns NULL, nil, undefined, or the closest equivalent in the corresponding language, if the specified key does not exist.
- set(*key, value [, expiry]*): Sets the item associated with a key in the cache to the specified value. This either updates an existing item if the key already exists, or adds a new key/value pair if

the key doesn't exist. If the expiry time is specified, then the item expires (and is deleted) when the expiry time is reached. The time is specified in seconds, and is taken as a relative time if the value is less than 30 days (30*24*60*60), or an absolute time (epoch) if larger than this value.

- `add(key, value [, expiry])`: Adds the key and associated value to the cache, if the specified key does not already exist.
- `replace(key, value [, expiry])`: Replaces the item associated with the specified `key`, only if the key already exists. The new value is given by the `value` parameter.
- `delete(key [, time])`: Deletes the `key` and its associated item from the cache. If you supply a `time`, then adding another item with the specified `key` is blocked for the specified period.
- `incr(key , value)`: Increments the item associated with the `key` by the specified `value`.
- `decr(key , value)`: Decrements the item associated with the `key` by the specified `value`.
- `flush_all`: Invalidates (or expires) all the current items in the cache. Technically they still exist (they are not deleted), but they are silently destroyed the next time you try to access them.

In all implementations, most or all of these functions are duplicated through the corresponding native language interface.

When practical, use `memcached` to store full items, rather than caching a single column value from the database. For example, when displaying a record about an object (invoice, user history, or blog post), load all the data for the associated entry from the database, and compile it into the internal structure that would normally be required by the application. Save the complete object in the cache.

Complex data structures cannot be stored directly. Most interfaces serialize the data for you, that is, put it in a textual form that can reconstruct the original pointers and nesting. Perl uses `Storable`, PHP uses `serialize`, Python uses `cPickle` (or `Pickle`) and Java uses the `Serializable` interface. In most cases, the serialization interface used is customizable. To share data stored in `memcached` instances between different language interfaces, consider using a common serialization solution such as JSON (Javascript Object Notation).

17.2.3.2 Using `memcached` as a MySQL Caching Layer

When using `memcached` to cache MySQL data, your application must retrieve data from the database and load the appropriate key-value pairs into the cache. Then, subsequent lookups can be done directly from the cache.

Because MySQL has its own in-memory caching mechanisms for queried data, such as the `InnoDB` buffer pool and the MySQL query cache, look for opportunities beyond loading individual column values or rows into the cache. Prefer to cache composite values, such as those retrieved from multiple tables through a join query, or result sets assembled from multiple rows.

Caution

Limit the information in the cache to non-sensitive data, because there is no security required to access or update the information within a `memcached` instance. Anybody with access to the machine has the ability to read, view and potentially update the information. To keep the data secure, encrypt the information before caching it. To restrict the users capable of connecting to the server, either disable network access, or use IPTables or similar techniques to restrict access to the `memcached` ports to a select set of hosts.

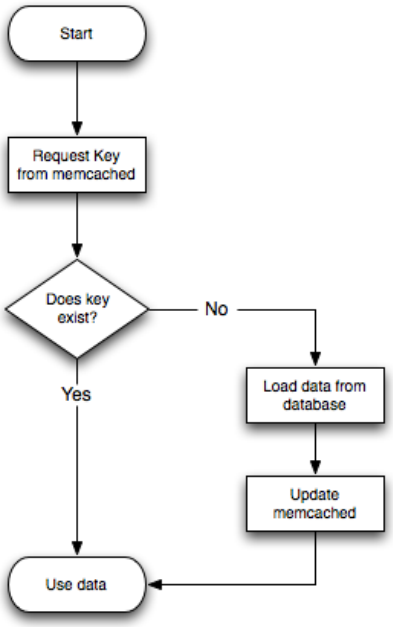
You can introduce `memcached` to an existing application, even if caching was not part of the original design. In many languages and environments the changes to the application will be just a few lines, first to attempt to read from the cache when loading data, fall back to the old method if the information is not cached, and to update the cache with information once the data has been read.

The general sequence for using memcached in any language as a caching solution for MySQL is as follows:

1. Request the item from the cache.
2. If the item exists, use the item data.
3. If the item does not exist, load the data from MySQL, and store the value into the cache. This means the value is available to the next client that requests it from the cache.

For a flow diagram of this sequence, see Figure 17.6, “Typical memcached Application Flowchart”.

Figure 17.6 Typical memcached Application Flowchart



Adapting Database Best Practices to memcached Applications

The most direct way to cache MySQL data is to use a 2-column table, where the first column is a primary key. Because of the uniqueness requirements for memcached keys, make sure your database schema makes appropriate use of primary keys and unique constraints.

If you combine multiple column values into a single memcached item value, choose data types to make it easy to parse the value back into its components, for example by using a separator character between numeric values.

The queries that map most easily to memcached lookups are those with a single WHERE clause, using an = or IN operator. For complicated WHERE clauses, or those using operators such as <, >, BETWEEN, or LIKE, memcached does not provide a simple or efficient way to scan through or filter the keys or associated values, so typically you perform those operations as SQL queries on the underlying database.

17.2.3.3 Using libmemcached with C and C++

The libmemcached library provides both C and C++ interfaces to memcached and is also the basis for a number of different additional API implementations, including Perl, Python and Ruby. Understanding the core libmemcached functions can help when using these other interfaces.

The C library is the most comprehensive interface library for memcached and provides functions and operational systems not always exposed in interfaces not based on the libmemcached library.

The different functions can be divided up according to their basic operation. In addition to functions that interface to the core API, a number of utility functions provide extended functionality, such as appending and prepending data.

To build and install libmemcached, download the libmemcached package, run configure, and then build and install:

```
shell> tar xjf libmemcached-0.21.tar.gz
shell> cd libmemcached-0.21
shell> ./configure
shell> make
shell> make install
```

On many Linux operating systems, you can install the corresponding libmemcached package through the usual yum, apt-get, or similar commands.

To build an application that uses the library, first set the list of servers. Either directly manipulate the servers configured within the main memcached_st structure, or separately populate a list of servers, and then add this list to the memcached_st structure. The latter method is used in the following example. Once the server list has been set, you can call the functions to store or retrieve data. A simple application for setting a preset value to localhost is provided here:

```
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#include <libmemcached/memcached.h>

int main(int argc, char *argv[])
{
  memcached_server_st *servers = NULL;
  memcached_st *memc;
  memcached_return rc;
  char *key= "keystring";
  char *value= "keyvalue";

  memcached_server_st *memcached_servers_parse (char *server_strings);
  memc= memcached_create(NULL);

  servers= memcached_server_list_append(servers, "localhost", 11211, &rc);
  rc= memcached_server_push(memc, servers);

  if (rc == MEMCACHED_SUCCESS)
    fprintf(stderr,"Added server successfully\n");
  else
    fprintf(stderr,"Couldn't add server: %s\n",memcached_strerror(memc, rc));

  rc= memcached_set(memc, key, strlen(key), value, strlen(value), (time_t)0, (uint32_t)0);

  if (rc == MEMCACHED_SUCCESS)
    fprintf(stderr,"Key stored successfully\n");
  else
    fprintf(stderr,"Couldn't store key: %s\n",memcached_strerror(memc, rc));

  return 0;
}
```

To test the success of an operation, use the return value, or populated result code, for a given function. The value is always set to MEMCACHED_SUCCESS if the operation succeeded. In the event of a failure, use the memcached_strerror() function to translate the result code into a printable string.

To build the application, specify the memcached library:

```
shell> gcc -o memc_basic memc_basic.c -lmemcached
```

Running the above sample application, after starting a memcached server, should return a success message:

```
shell> memc_basic
Added server successfully
Key stored successfully
```

libmemcached Base Functions

The base libmemcached functions let you create, destroy and clone the main memcached_st structure that is used to interface with the memcached servers. The main functions are defined below:

```
memcached_st *memcached_create (memcached_st *ptr);
```

Creates a new memcached_st structure for use with the other libmemcached API functions. You can supply an existing, static, memcached_st structure, or NULL to have a new structured allocated. Returns a pointer to the created structure, or NULL on failure.

```
void memcached_free (memcached_st *ptr);
```

Frees the structure and memory allocated to a previously created memcached_st structure.

```
memcached_st *memcached_clone(memcached_st *clone, memcached_st *source);
```

Clones an existing memcached structure from the specified source, copying the defaults and list of servers defined in the structure.

libmemcached Server Functions

The libmemcached API uses a list of servers, stored within the memcached_server_st structure, to act as the list of servers used by the rest of the functions. To use memcached, you first create the server list, and then apply the list of servers to a valid libmemcached object.

Because the list of servers, and the list of servers within an active libmemcached object can be manipulated separately, you can update and manage server lists while an active libmemcached interface is running.

The functions for manipulating the list of servers within a memcached_st structure are:

```
memcached_return
   memcached_server_add (memcached_st *ptr,
                         char *hostname,
                         unsigned int port);
```

Adds a server, using the given hostname and port into the memcached_st structure given in ptr.

```
memcached_return
   memcached_server_add_unix_socket (memcached_st *ptr,
                                     char *socket);
```

Adds a Unix socket to the list of servers configured in the memcached_st structure.

```
unsigned int memcached_server_count (memcached_st *ptr);
```

Returns a count of the number of configured servers within the memcached_st structure.

```
memcached_server_st *
   memcached_server_list (memcached_st *ptr);
```

Returns an array of all the defined hosts within a memcached_st structure.

```
memcached_return
   memcached_server_push (memcached_st *ptr,
                          memcached_server_st *list);
```

Pushes an existing list of servers onto list of servers configured for a current memcached_st structure. This adds servers to the end of the existing list, and duplicates are not checked.

The memcached_server_st structure can be used to create a list of memcached servers which can then be applied individually to memcached_st structures.

```
memcached_server_st *
   memcached_server_list_append (memcached_server_st *ptr,
                                 char *hostname,
                                 unsigned int port,
                                 memcached_return *error);
```

Adds a server, with hostname and port, to the server list in ptr. The result code is handled by the error argument, which should point to an existing memcached_return variable. The function returns a pointer to the returned list.

```
unsigned int memcached_server_list_count (memcached_server_st *ptr);
```

Returns the number of the servers in the server list.

```
void memcached_server_list_free (memcached_server_st *ptr);
```

Frees the memory associated with a server list.

```
memcached_server_st *memcached_servers_parse (char *server_strings);
```

Parses a string containing a list of servers, where individual servers are separated by a comma, space, or both, and where individual servers are of the form *server*[:*port*]. The return value is a server list structure.

libmemcached Set Functions

The set-related functions within libmemcached provide the same functionality as the core functions supported by the memcached protocol. The full definition for the different functions is the same for all the base functions (add, replace, prepend, append). For example, the function definition for memcached_set() is:

```
memcached_return
   memcached_set (memcached_st *ptr,
                  const char *key,
                  size_t key_length,
                  const char *value,
                  size_t value_length,
                  time_t expiration,
                  uint32_t flags);
```

The ptr is the memcached_st structure. The key and key_length define the key name and length, and value and value_length the corresponding value and length. You can also set the expiration and optional flags. For more information, see Controlling libmemcached Behaviors.

This table outlines the remainder of the set-related libmemcached functions and the equivalent core functions supported by the memcached protocol.

libmemcached Function	Equivalent Core Function
memcached_set(memc, key, key_length, value, value_length, expiration, flags)	Generic set() operation.
memcached_add(memc, key, key_length, value, value_length, expiration, flags)	Generic add() function.
memcached_replace(memc, key, key_length, value, value_length, expiration, flags)	Generic replace().
memcached_prepend(memc, key, key_length, value, value_length, expiration, flags)	Prepends the specified value before the current value of the specified key.
memcached_append(memc, key, key_length, value, value_length, expiration, flags)	Appends the specified value after the current value of the specified key.
memcached_cas(memc, key, key_length, value, value_length, expiration, flags, cas)	Overwrites the data for a given key as long as the corresponding cas value is still the same within the server.
memcached_set_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the generic set(), but has the option of an additional master key that can be used to identify an individual server.
memcached_add_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the generic add(), but has the option of an additional master key that can be used to identify an individual server.
memcached_replace_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the generic replace(), but has the option of an additional master key that can be used to identify an individual server.
memcached_prepend_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the memcached_prepend(), but has the option of an additional master key that can be used to identify an individual server.
memcached_append_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the memcached_append(), but has the option of an additional master key that can be used to identify an individual server.
memcached_cas_by_key(memc, master_key, master_key_length, key, key_length, value, value_length, expiration, flags)	Similar to the memcached_cas(), but has the option of an additional master key that can be used to identify an individual server.

The by_key methods add two further arguments that define the master key, to be used and applied during the hashing stage for selecting the servers. You can see this in the following definition:

```
memcached_return
   memcached_set_by_key(memcached_st *ptr,
                        const char *master_key,
                        size_t master_key_length,
                        const char *key,
                        size_t key_length,
                        const char *value,
```

```
                        size_t value_length,
                        time_t expiration,
                        uint32_t flags);
```

All the functions return a value of type memcached_return, which you can compare against the MEMCACHED_SUCCESS constant.

libmemcached Get Functions

The libmemcached functions provide both direct access to a single item, and a multiple-key request mechanism that provides much faster responses when fetching a large number of keys simultaneously.

The main get-style function, which is equivalent to the generic get() is memcached_get(). This function returns a string pointer, pointing to the value associated with the specified key.

```
char *memcached_get (memcached_st *ptr,
                     const char *key, size_t key_length,
                     size_t *value_length,
                     uint32_t *flags,
                     memcached_return *error);
```

A multi-key get, memcached_mget(), is also available. Using a multiple key get operation is much quicker to do in one block than retrieving the key values with individual calls to memcached_get(). To start the multi-key get, call memcached_mget():

```
memcached_return
    memcached_mget (memcached_st *ptr,
                    char **keys, size_t *key_length,
                    unsigned int number_of_keys);
```

The return value is the success of the operation. The keys parameter should be an array of strings containing the keys, and key_length an array containing the length of each corresponding key. number_of_keys is the number of keys supplied in the array.

To fetch the individual values, use memcached_fetch() to get each corresponding value.

```
char *memcached_fetch (memcached_st *ptr,
                       const char *key, size_t *key_length,
                       size_t *value_length,
                       uint32_t *flags,
                       memcached_return *error);
```

The function returns the key value, with the key, key_length and value_length parameters being populated with the corresponding key and length information. The function returns NULL when there are no more values to be returned. A full example, including the populating of the key data and the return of the information is provided here.

```
#include <stdio.h>
#include <sstring.h>
#include <unistd.h>
#include <libmemcached/memcached.h>

int main(int argc, char *argv[])
{
  memcached_server_st *servers = NULL;
  memcached_st *memc;
  memcached_return rc;
  char *keys[]= {"huey", "dewey", "louie"};
  size_t key_length[3];
  char *values[]= {"red", "blue", "green"};
  size_t value_length[3];
  unsigned int x;
  uint32_t flags;
```

```
  char return_key[MEMCACHED_MAX_KEY];
  size_t return_key_length;
  char *return_value;
  size_t return_value_length;

  memc= memcached_create(NULL);

  servers= memcached_server_list_append(servers, "localhost", 11211, &rc);
  rc= memcached_server_push(memc, servers);

  if (rc == MEMCACHED_SUCCESS)
    fprintf(stderr,"Added server successfully\n");
  else
    fprintf(stderr,"Couldn't add server: %s\n",memcached_strerror(memc, rc));

  for(x= 0; x < 3; x++)
    {
      key_length[x] = strlen(keys[x]);
      value_length[x] = strlen(values[x]);

      rc= memcached_set(memc, keys[x], key_length[x], values[x],
                        value_length[x], (time_t)0, (uint32_t)0);
      if (rc == MEMCACHED_SUCCESS)
        fprintf(stderr,"Key %s stored successfully\n",keys[x]);
      else
        fprintf(stderr,"Couldn't store key: %s\n",memcached_strerror(memc, rc));
    }

  rc= memcached_mget(memc, keys, key_length, 3);

  if (rc == MEMCACHED_SUCCESS)
    {
      while ((return_value= memcached_fetch(memc, return_key, &return_key_length,
                                            &return_value_length, &flags, &rc)) != NULL)
        {
          if (rc == MEMCACHED_SUCCESS)
            {
              fprintf(stderr,"Key %s returned %s\n",return_key, return_value);
            }
        }
    }

  return 0;
}
```

Running the above application produces the following output:

```
shell> memc_multi_fetch
Added server successfully
Key huey stored successfully
Key dewey stored successfully
Key louie stored successfully
Key huey returned red
Key dewey returned blue
Key louie returned green
```

Controlling libmemcached Behaviors

The behavior of libmemcached can be modified by setting one or more behavior flags. These can either be set globally, or they can be applied during the call to individual functions. Some behaviors also accept an additional setting, such as the hashing mechanism used when selecting servers.

To set global behaviors:

```
memcached_return
   memcached_behavior_set (memcached_st *ptr,
                           memcached_behavior flag,
                           uint64_t data);
```

To get the current behavior setting:

```
uint64_t
   memcached_behavior_get (memcached_st *ptr,
                           memcached_behavior flag);
```

The following table describes libmemcached behavior flags.

Behavior	Description
MEMCACHED_BEHAVIOR_NO_BLOCK	Caused libmemcached to use asynchronous I/O.
MEMCACHED_BEHAVIOR_TCP_NODELAY	Turns on no-delay for network sockets.
MEMCACHED_BEHAVIOR_HASH	Without a value, sets the default hashing algorithm for keys to use MD5. Other valid values include MEMCACHED_HASH_DEFAULT, MEMCACHED_HASH_MD5, MEMCACHED_HASH_CRC, MEMCACHED_HASH_FNV1_64, MEMCACHED_HASH_FNV1A_64, MEMCACHED_HASH_FNV1_32, and MEMCACHED_HASH_FNV1A_32.
MEMCACHED_BEHAVIOR_DISTRIBUTION	Changes the method of selecting the server used to store a given value. The default method is MEMCACHED_DISTRIBUTION_MODULA. You can enable consistent hashing by setting MEMCACHED_DISTRIBUTION_CONSISTENT. MEMCACHED_DISTRIBUTION_CONSISTENT is an alias for the value MEMCACHED_DISTRIBUTION_CONSISTENT_KETAMA.
MEMCACHED_BEHAVIOR_CACHE_LOOKUPS	Cache the lookups made to the DNS service. This can improve the performance if you are using names instead of IP addresses for individual hosts.
MEMCACHED_BEHAVIOR_SUPPORT_CAS	Support CAS operations. By default, this is disabled because it imposes a performance penalty.
MEMCACHED_BEHAVIOR_KETAMA	Sets the default distribution to MEMCACHED_DISTRIBUTION_CONSISTENT_KETAMA and the hash to MEMCACHED_HASH_MD5.
MEMCACHED_BEHAVIOR_POLL_TIMEOUT	Modify the timeout value used by poll(). Supply a signed int pointer for the timeout value.
MEMCACHED_BEHAVIOR_BUFFER_REQUESTS	Buffers IO requests instead of them being sent. A get operation, or closing the connection causes the data to be flushed.
MEMCACHED_BEHAVIOR_VERIFY_KEY	Forces libmemcached to verify that a specified key is valid.
MEMCACHED_BEHAVIOR_SORT_HOSTS	If set, hosts added to the list of configured hosts for a memcached_st structure are placed into the host list in sorted order. This breaks consistent hashing if that behavior has been enabled.
MEMCACHED_BEHAVIOR_CONNECT_TIMEOUT	In nonblocking mode this changes the value of the timeout during socket connection.

libmemcached Command-Line Utilities

In addition to the main C library interface, libmemcached also includes a number of command-line utilities that can be useful when working with and debugging memcached applications.

All of the command-line tools accept a number of arguments, the most critical of which is servers, which specifies the list of servers to connect to when returning information.

The main tools are:

- memcat: Display the value for each ID given on the command line:

```
shell> memcat --servers=localhost hwkey
Hello world
```

- memcp: Copy the contents of a file into the cache, using the file name as the key:

```
shell> echo "Hello World" > hwkey
shell> memcp --servers=localhost hwkey
shell> memcat --servers=localhost hwkey
Hello world
```

- memrm: Remove an item from the cache:

```
shell> memcat --servers=localhost hwkey
Hello world
shell> memrm --servers=localhost hwkey
shell> memcat --servers=localhost hwkey
```

- memslap: Test the load on one or more memcached servers, simulating get/set and multiple client operations. For example, you can simulate the load of 100 clients performing get operations:

```
shell> memslap --servers=localhost --concurrency=100 --flush --test=get
memslap --servers=localhost --concurrency=100 --flush --test=get Threads connecting to servers 100
 Took 13.571 seconds to read data
```

- memflush: Flush (empty) the contents of the memcached cache.

```
shell> memflush --servers=localhost
```

17.2.3.4 Using MySQL and memcached with Perl

The Cache::Memcached module provides a native interface to the Memcache protocol, and provides support for the core functions offered by memcached. Install the module using your operating system's package management system, or using CPAN:

```
root-shell> perl -MCPAN -e 'install Cache::Memcached'
```

To use memcached from Perl through the Cache::Memcached module, first create a new Cache::Memcached object that defines the list of servers and other parameters for the connection. The only argument is a hash containing the options for the cache interface. For example, to create a new instance that uses three memcached servers:

```
use Cache::Memcached;

my $cache = new Cache::Memcached {
    'servers' => [
        '192.168.0.100:11211',
        '192.168.0.101:11211',
        '192.168.0.102:11211',
 ],
};
```

Note

When using the Cache::Memcached interface with multiple servers, the API automatically performs certain operations across all the servers in the group. For example, getting statistical information through Cache::Memcached

> returns a hash that contains data on a host-by-host basis, as well as generalized statistics for all the servers in the group.

You can set additional properties on the cache object instance when it is created by specifying the option as part of the option hash. Alternatively, you can use a corresponding method on the instance:

- `servers` or method `set_servers()`: Specifies the list of the servers to be used. The servers list should be a reference to an array of servers, with each element as the address and port number combination (separated by a colon). You can also specify a local connection through a Unix socket (for example `/tmp/sock/memcached`). To specify the server with a weight (indicating how much more frequently the server should be used during hashing), specify an array reference with the `memcached` server instance and a weight number. Higher numbers give higher priority.
- `compress_threshold` or method `set_compress_threshold()`: Specifies the threshold when values are compressed. Values larger than the specified number are automatically compressed (using `zlib`) during storage and retrieval.
- `no_rehash` or method `set_norehash()`: Disables finding a new server if the original choice is unavailable.
- `readonly` or method `set_readonly()`: Disables writes to the `memcached` servers.

Once the `Cache::Memcached` object instance has been configured, you can use the `set()` and `get()` methods to store and retrieve information from the `memcached` servers. Objects stored in the cache are automatically serialized and deserialized using the `Storable` module.

The `Cache::Memcached` interface supports the following methods for storing/retrieving data, and relate to the generic methods as shown in the table.

`Cache::Memcached` Function	**Equivalent Generic Method**
`get()`	Generic `get()`.
`get_multi(keys)`	Gets multiple `keys` from memcache using just one query. Returns a hash reference of key/value pairs.
`set()`	Generic `set()`.
`add()`	Generic `add()`.
`replace()`	Generic `replace()`.
`delete()`	Generic `delete()`.
`incr()`	Generic `incr()`.
`decr()`	Generic `decr()`.

Below is a complete example for using `memcached` with Perl and the `Cache::Memcached` module:

```
#!/usr/bin/perl

use Cache::Memcached;
use DBI;
use Data::Dumper;

# Configure the memcached server

my $cache = new Cache::Memcached {
    'servers' => [
                   'localhost:11211',
                   ],
    };
```

```
# Get the film name from the command line
# memcached keys must not contain spaces, so create
# a key name by replacing spaces with underscores

my $filmname = shift or die "Must specify the film name\n";
my $filmkey = $filmname;
$filmkey =~ s/ /_/;

# Load the data from the cache

my $filmdata = $cache->get($filmkey);

# If the data wasn't in the cache, then we load it from the database

if (!defined($filmdata))
{
    $filmdata = load_filmdata($filmname);

    if (defined($filmdata))
    {

# Set the data into the cache, using the key

 if ($cache->set($filmkey,$filmdata))
        {
            print STDERR "Film data loaded from database and cached\n";
        }
        else
        {
            print STDERR "Couldn't store to cache\n";
 }
    }
    else
    {
      die "Couldn't find $filmname\n";
    }
}
else
{
    print STDERR "Film data loaded from Memcached\n";
}

sub load_filmdata
{
    my ($filmname) = @_;

    my $dsn = "DBI:mysql:database=sakila;host=localhost;port=3306";

    $dbh = DBI->connect($dsn, 'sakila','password');

    my ($filmbase) = $dbh->selectrow_hashref(sprintf('select * from film where title = %s',
                                                     $dbh->quote($filmname)));

    if (!defined($filmname))
    {
      return (undef);
    }

    $filmbase->{stars} =
 $dbh->selectall_arrayref(sprintf('select concat(first_name," ",last_name) ' .
                                        'from film_actor left join (actor) ' .
                                        'on (film_actor.actor_id = actor.actor_id) ' .
                                        ' where film_id=%s',
                                        $dbh->quote($filmbase->{film_id})));

    return($filmbase);
}
```

The example uses the Sakila database, obtaining film data from the database and writing a composite record of the film and actors to memcached. When calling it for a film does not exist, you get this result:

```
shell> memcached-sakila.pl "ROCK INSTINCT"
Film data loaded from database and cached
```

When accessing a film that has already been added to the cache:

```
shell> memcached-sakila.pl "ROCK INSTINCT"
Film data loaded from Memcached
```

17.2.3.5 Using MySQL and memcached with Python

The Python memcache module interfaces to memcached servers, and is written in pure Python (that is, without using one of the C APIs). You can download and install a copy from Python Memcached.

To install, download the package and then run the Python installer:

```
python setup.py install
running install
running bdist_egg
running egg_info
creating python_memcached.egg-info
...
removing 'build/bdist.linux-x86_64/egg' (and everything under it)
Processing python_memcached-1.43-py2.4.egg
creating /usr/lib64/python2.4/site-packages/python_memcached-1.43-py2.4.egg
Extracting python_memcached-1.43-py2.4.egg to /usr/lib64/python2.4/site-packages
Adding python-memcached 1.43 to easy-install.pth file

Installed /usr/lib64/python2.4/site-packages/python_memcached-1.43-py2.4.egg
Processing dependencies for python-memcached==1.43
Finished processing dependencies for python-memcached==1.43
```

Once installed, the memcache module provides a class-based interface to your memcached servers. When you store Python data structures as memcached items, they are automatically serialized (turned into string values) using the Python cPickle or pickle modules.

To create a new memcache interface, import the memcache module and create a new instance of the memcache.Client class. For example, if the memcached daemon is running on localhost using the default port:

```
import memcache
memc = memcache.Client(['127.0.0.1:11211'])
```

The first argument is an array of strings containing the server and port number for each memcached instance to use. To enable debugging, set the optional debug parameter to 1.

By default, the hashing mechanism used to divide the items among multiple servers is crc32. To change the function used, set the value of memcache.serverHashFunction to the alternate function to use. For example:

```
from zlib import adler32
memcache.serverHashFunction = adler32
```

Once you have defined the servers to use within the memcache instance, the core functions provide the same functionality as in the generic interface specification. The following table provides a summary of the supported functions:

Python memcache Function	Equivalent Generic Function
get()	Generic get().

Python memcache Function	Equivalent Generic Function
get_multi(keys)	Gets multiple values from the supplied array of keys. Returns a hash reference of key/value pairs.
set()	Generic set().
set_multi(dict [, expiry [, key_prefix]])	Sets multiple key/value pairs from the supplied dict.
add()	Generic add().
replace()	Generic replace().
prepend(key, value [, expiry])	Prepends the supplied value to the value of the existing key.
append(key, value [, expiry[)	Appends the supplied value to the value of the existing key.
delete()	Generic delete().
delete_multi(keys [, expiry [, key_prefix]])	Deletes all the keys from the hash matching each string in the array keys.
incr()	Generic incr().
decr()	Generic decr().

Note

Within the Python memcache module, all the *_multi()functions support an optional key_prefix parameter. If supplied, then the string is used as a prefix to all key lookups. For example, if you call:

```
memc.get_multi(['a','b'], key_prefix='users:')
```

The function retrieves the keys users:a and users:b from the servers.

Here is an example showing the storage and retrieval of information to a memcache instance, loading the raw data from MySQL:

```
import sys
import MySQLdb
import memcache

memc = memcache.Client(['127.0.0.1:11211'], debug=1);

try:
    conn = MySQLdb.connect (host = "localhost",
                            user = "sakila",
                            passwd = "password",
                            db = "sakila")
except MySQLdb.Error, e:
     print "Error %d: %s" % (e.args[0], e.args[1])
     sys.exit (1)

popularfilms = memc.get('top5films')

if not popularfilms:
    cursor = conn.cursor()
    cursor.execute('select film_id,title from film order by rental_rate desc limit 5')
    rows = cursor.fetchall()
    memc.set('top5films',rows,60)
    print "Updated memcached with MySQL data"
else:
    print "Loaded data from memcached"
    for row in popularfilms:
        print "%s, %s" % (row[0], row[1])
```

When executed for the first time, the data is loaded from the MySQL database and stored to the memcached server.

```
shell> python memc_python.py
Updated memcached with MySQL data
```

Because the data is automatically serialized using cPickle/pickle, when you load the data back from memcached, you can use the object directly. In the example above, the information stored to memcached is in the form of rows from a Python DB cursor. When accessing the information (within the 60 second expiry time), the data is loaded from memcached and dumped:

```
shell> python memc_python.py
Loaded data from memcached
2, ACE GOLDFINGER
7, AIRPLANE SIERRA
8, AIRPORT POLLOCK
10, ALADDIN CALENDAR
13, ALI FOREVER
```

The serialization and deserialization happens automatically. Because serialization of Python data may be incompatible with other interfaces and languages, you can change the serialization module used during initialization. For example, you might use JSON format when you store complex data structures using a script written in one language, and access them in a script written in a different language.

17.2.3.6 Using MySQL and memcached with PHP

PHP provides support for the Memcache functions through a PECL extension. To enable the PHP memcache extensions, build PHP using the --enable-memcache option to configure when building from source.

If you are installing on a Red Hat-based server, you can install the php-pecl-memcache RPM:

```
root-shell> yum --install php-pecl-memcache
```

On Debian-based distributions, use the php-memcache package.

To set global runtime configuration options, specify the configuration option values within your php.ini file. The following table provides the name, default value, and a description for each global runtime configuration option.

Configuration option	Default	Description
memcache.allow_failover	1	Specifies whether another server in the list should be queried if the first server selected fails.
memcache.max_failover_attempts	20	Specifies the number of servers to try before returning a failure.
memcache.chunk_size	8192	Defines the size of network chunks used to exchange data with the memcached server.
memcache.default_port	11211	Defines the default port to use when communicating with the memcached servers.
memcache.hash_strategy	standard	Specifies which hash strategy to use. Set to consistent to enable servers to be added or removed from the pool without causing the keys to be remapped to other servers. When set to standard, an older (modula) strategy

Configuration option	Default	Description
		is used that potentially uses different servers for storage.
memcache.hash_function	crc32	Specifies which function to use when mapping keys to servers. crc32 uses the standard CRC32 hash. fnv uses the FNV-1a hashing algorithm.

To create a connection to a memcached server, create a new Memcache object and then specify the connection options. For example:

```
<?php

$cache = new Memcache;
$cache->connect('localhost',11211);
?>
```

This opens an immediate connection to the specified server.

To use multiple memcached servers, you need to add servers to the memcache object using addServer():

```
bool Memcache::addServer ( string $host [, int $port [, bool $persistent
                [, int $weight [, int $timeout [, int $retry_interval
                [, bool $status [, callback $failure_callback
                ]]]]]] )
```

The server management mechanism within the php-memcache module is a critical part of the interface as it controls the main interface to the memcached instances and how the different instances are selected through the hashing mechanism.

To create a simple connection to two memcached instances:

```
<?php

$cache = new Memcache;
$cache->addServer('192.168.0.100',11211);
$cache->addServer('192.168.0.101',11211);
?>
```

In this scenario, the instance connection is not explicitly opened, but only opened when you try to store or retrieve a value. To enable persistent connections to memcached instances, set the $persistent argument to true. This is the default setting, and causes the connections to remain open.

To help control the distribution of keys to different instances, use the global memcache.hash_strategy setting. This sets the hashing mechanism used to select. You can also add another weight to each server, which effectively increases the number of times the instance entry appears in the instance list, therefore increasing the likelihood of the instance being chosen over other instances. To set the weight, set the value of the $weight argument to more than one.

The functions for setting and retrieving information are identical to the generic functional interface offered by memcached, as shown in this table:

PECL memcache Function	Generic Function
get()	Generic get().
set()	Generic set().
add()	Generic add().

PECL memcache Function	Generic Function
replace()	Generic replace().
delete()	Generic delete().
increment()	Generic incr().
decrement()	Generic decr().

A full example of the PECL memcache interface is provided below. The code loads film data from the Sakila database when the user provides a film name. The data stored into the memcached instance is recorded as a mysqli result row, and the API automatically serializes the information for you.

```
<?php

$memc = new Memcache;
$memc->addServer('localhost','11211');

if(empty($_POST['film'])) {
?>
  <html xmlns="http://www.w3.org/1999/xhtml" xml:lang="en" lang="en">
    <head>
      <meta http-equiv="Content-Type" content="text/html; charset=utf-8" />
      <title>Simple Memcache Lookup</title>
    </head>
    <body>
      <form method="post">
        <p><b>Film</b>: <input type="text" size="20" name="film"></p>
        <input type="submit">
      </form>
      <hr/>
<?php

} else {

    echo "Loading data...\n";

    $film   = htmlspecialchars($_POST['film'], ENT_QUOTES, 'UTF-8');
    $mfilms = $memc->get($film);

    if ($mfilms) {

        printf("<p>Film data for %s loaded from memcache</p>", $mfilms['title']);

        foreach (array_keys($mfilms) as $key) {
            printf("<p><b>%s</b>: %s</p>", $key, $mfilms[$key]);
        }

    } else {

        $mysqli = mysqli('localhost','sakila','password','sakila');

        if (mysqli_connect_error()) {
            sprintf("Database error: (%d) %s", mysqli_connect_errno(), mysqli_connect_error());
            exit;
        }

        $sql = sprintf('SELECT * FROM film WHERE title="%s"', $mysqli->real_escape_string($film));

        $result = $mysqli->query($sql);

        if (!$result) {
            sprintf("Database error: (%d) %s", $mysqli->errno, $mysqli->error);
            exit;
        }

        $row = $result->fetch_assoc();

        $memc->set($row['title'], $row);
```

```
        printf("<p>Loaded (%s) from MySQL</p>", htmlspecialchars($row['title'], ENT_QUOTES, 'UTF-8');
    }
}
?>
  </body>
</html>
```

With PHP, the connections to the memcached instances are kept open as long as the PHP and associated Apache instance remain running. When adding or removing servers from the list in a running instance (for example, when starting another script that mentions additional servers), the connections are shared, but the script only selects among the instances explicitly configured within the script.

To ensure that changes to the server list within a script do not cause problems, make sure to use the consistent hashing mechanism.

17.2.3.7 Using MySQL and memcached with Ruby

There are a number of different modules for interfacing to memcached within Ruby. The Ruby-MemCache client library provides a native interface to memcached that does not require any external libraries, such as libmemcached. You can obtain the installer package from http://www.deveiate.org/projects/RMemCache.

To install, extract the package and then run install.rb:

```
shell> install.rb
```

If you have RubyGems, you can install the Ruby-MemCache gem:

```
shell> gem install Ruby-MemCache
Bulk updating Gem source index for: http://gems.rubyforge.org
Install required dependency io-reactor? [Yn]  y
Successfully installed Ruby-MemCache-0.0.1
Successfully installed io-reactor-0.05
Installing ri documentation for io-reactor-0.05...
Installing RDoc documentation for io-reactor-0.05...
```

To use a memcached instance from within Ruby, create a new instance of the MemCache object.

```
require 'memcache'
memc = MemCache::new '192.168.0.100:11211'
```

You can add a weight to each server to increase the likelihood of the server being selected during hashing by appending the weight count to the server host name/port string:

```
require 'memcache'
memc = MemCache::new '192.168.0.100:11211:3'
```

To add servers to an existing list, you can append them directly to the MemCache object:

```
memc += ["192.168.0.101:11211"]
```

To set data into the cache, you can just assign a value to a key within the new cache object, which works just like a standard Ruby hash object:

```
memc["key"] = "value"
```

Or to retrieve the value:

```
print memc["key"]
```

For more explicit actions, you can use the method interface, which mimics the main memcached API functions, as summarized in the following table:

Ruby MemCache Method	Equivalent memcached API Functions
get()	Generic get().
get_hash(keys)	Get the values of multiple keys, returning the information as a hash of the keys and their values.
set()	Generic set().
set_many(pairs)	Set the values of the keys and values in the hash pairs.
add()	Generic add().
replace()	Generic replace().
delete()	Generic delete().
incr()	Generic incr().
decr()	Generic decr().

17.2.3.8 Using MySQL and memcached with Java

The com.danga.MemCached class within Java provides a native interface to memcached instances. You can obtain the client from https://github.com/gwhalin/Memcached-Java-Client/downloads. The Java class uses hashes that are compatible with libmemcached, so you can mix and match Java and libmemcached applications accessing the same memcached instances. The serialization between Java and other interfaces are not compatible. If this is a problem, use JSON or a similar nonbinary serialization format.

On most systems, you can download the package and use the jar directly.

To use the com.danga.MemCached interface, you create a MemCachedClient instance and then configure the list of servers by configuring the SockIOPool. Through the pool specification you set up the server list, weighting, and the connection parameters to optimized the connections between your client and the memcached instances that you configure.

Generally, you can configure the memcached interface once within a single class, then use this interface throughout the rest of your application.

For example, to create a basic interface, first configure the MemCachedClient and base SockIOPool settings:

```
public class MyClass {

    protected static MemCachedClient mcc = new MemCachedClient();

    static {

        String[] servers =
            {
                "localhost:11211",
            };

        Integer[] weights = { 1 };

        SockIOPool pool = SockIOPool.getInstance();
```

```
        pool.setServers( servers );
        pool.setWeights( weights );
```

In the above sample, the list of servers is configured by creating an array of the memcached instances to use. You can then configure individual weights for each server.

The remainder of the properties for the connection are optional, but you can set the connection numbers (initial connections, minimum connections, maximum connections, and the idle timeout) by setting the pool parameters:

```
pool.setInitConn( 5 );
pool.setMinConn( 5 );
pool.setMaxConn( 250 );
pool.setMaxIdle( 1000 * 60 * 60 * 6
```

Once the parameters have been configured, initialize the connection pool:

```
pool.initialize();
```

The pool, and the connection to your memcached instances should now be ready to use.

To set the hashing algorithm used to select the server used when storing a given key, use pool.setHashingAlg():

```
pool.setHashingAlg( SockIOPool.NEW_COMPAT_HASH );
```

Valid values are NEW_COMPAT_HASH, OLD_COMPAT_HASH and NATIVE_HASH are also basic modula hashing algorithms. For a consistent hashing algorithm, use CONSISTENT_HASH. These constants are equivalent to the corresponding hash settings within libmemcached.

The following table outlines the Java com.danga.MemCached methods and the equivalent generic methods in the memcached interface specification.

Java com.danga.MemCached Method	Equivalent Generic Method
get()	Generic get().
getMulti(keys)	Get the values of multiple keys, returning the information as Hash map using java.lang.String for the keys and java.lang.Object for the corresponding values.
set()	Generic set().
add()	Generic add().
replace()	Generic replace().
delete()	Generic delete().
incr()	Generic incr().
decr()	Generic decr().

17.2.3.9 Using the memcached TCP Text Protocol

Communicating with a memcached server can be achieved through either the TCP or UDP protocols. When using the TCP protocol, you can use a simple text based interface for the exchange of information.

When communicating with memcached, you can connect to the server using the port configured for the server. You can open a connection with the server without requiring authorization or login. As soon as you have connected, you can start to send commands to the server. When you have finished, you can terminate the connection without sending any specific disconnection command. Clients are encouraged to keep their connections open to decrease latency and improve performance.

Data is sent to the memcached server in two forms:

- Text lines, which are used to send commands to the server, and receive responses from the server.
- Unstructured data, which is used to receive or send the value information for a given key. Data is returned to the client in exactly the format it was provided.

Both text lines (commands and responses) and unstructured data are always terminated with the string \r\n. Because the data being stored may contain this sequence, the length of the data (returned by the client before the unstructured data is transmitted should be used to determine the end of the data.

Commands to the server are structured according to their operation:

- **Storage commands**: set, add, replace, append, prepend, cas

  Storage commands to the server take the form:

  ```
  command key [flags] [exptime] length [noreply]
  ```

  Or when using compare and swap (cas):

  ```
  cas key [flags] [exptime] length [casunique] [noreply]
  ```

  Where:

  - command: The command name.
    - set: Store value against key
    - add: Store this value against key if the key does not already exist
    - replace: Store this value against key if the key already exists
    - append: Append the supplied value to the end of the value for the specified key. The flags and exptime arguments should not be used.
    - prepend: Append value currently in the cache to the end of the supplied value for the specified key. The flags and exptime arguments should not be used.
    - cas: Set the specified key to the supplied value, only if the supplied casunique matches. This is effectively the equivalent of change the information if nobody has updated it since I last fetched it.
  - key: The key. All data is stored using a the specific key. The key cannot contain control characters or whitespace, and can be up to 250 characters in size.
  - flags: The flags for the operation (as an integer). Flags in memcached are transparent. The memcached server ignores the contents of the flags. They can be used by the client to indicate any type of information. In memcached 1.2.0 and lower the value is a 16-bit integer value. In memcached 1.2.1 and higher the value is a 32-bit integer.
  - exptime: The expiry time, or zero for no expiry.
  - length: The length of the supplied value block in bytes, excluding the terminating \r\n characters.

- casunique: A unique 64-bit value of an existing entry. This is used to compare against the existing value. Use the value returned by the gets command when issuing cas updates.
- noreply: Tells the server not to reply to the command.

For example, to store the value abcdef into the key xyzkey, you would use:

```
set xyzkey 0 0 6\r\nabcdef\r\n
```

The return value from the server is one line, specifying the status or error information. For more information, see Table 17.3, "memcached Protocol Responses".

- **Retrieval commands**: get, gets

Retrieval commands take the form:

```
get key1 [key2 .... keyn]
gets key1 [key2 ... keyn]
```

You can supply multiple keys to the commands, with each requested key separated by whitespace.

The server responds with an information line of the form:

```
VALUE key flags bytes [casunique]
```

Where:

- key: The key name.
- flags: The value of the flag integer supplied to the memcached server when the value was stored.
- bytes: The size (excluding the terminating \r\n character sequence) of the stored value.
- casunique: The unique 64-bit integer that identifies the item.

The information line is immediately followed by the value data block. For example:

```
get xyzkey\r\n
VALUE xyzkey 0 6\r\n
abcdef\r\n
```

If you have requested multiple keys, an information line and data block is returned for each key found. If a requested key does not exist in the cache, no information is returned.

- **Delete commands**: delete

Deletion commands take the form:

```
delete key [time] [noreply]
```

Where:

- key: The key name.
- time: The time in seconds (or a specific Unix time) for which the client wishes the server to refuse add or replace commands on this key. All add, replace, get, and gets commands fail during this period. set operations succeed. After this period, the key is deleted permanently and all commands are accepted.

If not supplied, the value is assumed to be zero (delete immediately).

- noreply: Tells the server not to reply to the command.

Responses to the command are either DELETED to indicate that the key was successfully removed, or NOT_FOUND to indicate that the specified key could not be found.

- **Increment/Decrement**: incr, decr

  The increment and decrement commands change the value of a key within the server without performing a separate get/set sequence. The operations assume that the currently stored value is a 64-bit integer. If the stored value is not a 64-bit integer, then the value is assumed to be zero before the increment or decrement operation is applied.

  Increment and decrement commands take the form:

```
incr key value [noreply]
decr key value [noreply]
```

  Where:

  - key: The key name.
  - value: An integer to be used as the increment or decrement value.
  - noreply: Tells the server not to reply to the command.

  The response is:

  - NOT_FOUND: The specified key could not be located.
  - value: The new value associated with the specified key.

  Values are assumed to be unsigned. For decr operations, the value is never decremented below 0. For incr operations, the value wraps around the 64-bit maximum.

- **Statistics commands**: stats

  The stats command provides detailed statistical information about the current status of the memcached instance and the data it is storing.

  Statistics commands take the form:

```
STAT [name] [value]
```

  Where:

  - name: The optional name of the statistics to return. If not specified, the general statistics are returned.
  - value: A specific value to be used when performing certain statistics operations.

  The return value is a list of statistics data, formatted as follows:

```
STAT name value
```

  The statistics are terminated with a single line, END.

  For more information, see Section 17.2.4, “Getting memcached Statistics”.

For reference, a list of the different commands supported and their formats is provided below.

Table 17.2 memcached Command Reference

Command	Command Formats
set	set key flags exptime length, set key flags exptime length noreply
add	add key flags exptime length, add key flags exptime length noreply
replace	replace key flags exptime length, replace key flags exptime length noreply
append	append key length, append key length noreply
prepend	prepend key length, prepend key length noreply
cas	cas key flags exptime length casunique, cas key flags exptime length casunique noreply
get	get key1 [key2 ... keyn]
gets	
delete	delete key, delete key noreply, delete key expiry, delete key expiry noreply
incr	incr key, incr key noreply, incr key value, incr key value noreply
decr	decr key, decr key noreply, decr key value, decr key value noreply
stat	stat, stat name, stat name value

When sending a command to the server, the response from the server is one of the settings in the following table. All response values from the server are terminated by \r\n:

Table 17.3 memcached Protocol Responses

String	Description
STORED	Value has successfully been stored.
NOT_STORED	The value was not stored, but not because of an error. For commands where you are adding a or updating a value if it exists (such as add and replace), or where the item has already been set to be deleted.
EXISTS	When using a cas command, the item you are trying to store already exists and has been modified since you last checked it.
NOT_FOUND	The item you are trying to store, update or delete does not exist or has already been deleted.
ERROR	You submitted a nonexistent command name.
CLIENT_ERROR errorstring	There was an error in the input line, the detail is contained in errorstring.
SERVER_ERROR errorstring	There was an error in the server that prevents it from returning the information. In extreme conditions, the server may disconnect the client after this error occurs.
VALUE keys flags length	The requested key has been found, and the stored key, flags and data block are returned, of the specified length.
DELETED	The requested key was deleted from the server.
STAT name value	A line of statistics data.

String	Description
END	The end of the statistics data.

17.2.4 Getting memcached Statistics

The memcached system has a built-in statistics system that collects information about the data being stored into the cache, cache hit ratios, and detailed information on the memory usage and distribution of information through the slab allocation used to store individual items. Statistics are provided at both a basic level that provide the core statistics, and more specific statistics for specific areas of the memcached server.

This information can be useful to ensure that you are getting the correct level of cache and memory usage, and that your slab allocation and configuration properties are set at an optimal level.

The stats interface is available through the standard memcached protocol, so the reports can be accessed by using telnet to connect to the memcached. The supplied memcached-tool includes support for obtaining the Section 17.2.4.2, “memcached Slabs Statistics” and Section 17.2.4.1, “memcached General Statistics” information. For more information, see Section 17.2.4.6, “Using memcached-tool”.

Alternatively, most of the language API interfaces provide a function for obtaining the statistics from the server.

For example, to get the basic stats using telnet:

```
shell> telnet localhost 11211
Trying ::1...
Connected to localhost.
Escape character is '^]'.
stats
STAT pid 23599
STAT uptime 675
STAT time 1211439587
STAT version 1.2.5
STAT pointer_size 32
STAT rusage_user 1.404992
STAT rusage_system 4.694685
STAT curr_items 32
STAT total_items 56361
STAT bytes 2642
STAT curr_connections 53
STAT total_connections 438
STAT connection_structures 55
STAT cmd_get 113482
STAT cmd_set 80519
STAT get_hits 78926
STAT get_misses 34556
STAT evictions 0
STAT bytes_read 6379783
STAT bytes_written 4860179
STAT limit_maxbytes 67108864
STAT threads 1
END
```

When using Perl and the Cache::Memcached module, the stats() function returns information about all the servers currently configured in the connection object, and total statistics for all the memcached servers as a whole.

For example, the following Perl script obtains the stats and dumps the hash reference that is returned:

```
use Cache::Memcached;
use Data::Dumper;

my $memc = new Cache::Memcached;
```

```
$memc->set_servers(\@ARGV);

print Dumper($memc->stats());
```

When executed on the same memcached as used in the Telnet example above we get a hash reference with the host by host and total statistics:

```
$VAR1 = {
    'hosts' => {
           'localhost:11211' => {
                      'misc' => {
                            'bytes' => '2421',
                            'curr_connections' => '3',
                            'connection_structures' => '56',
                            'pointer_size' => '32',
                            'time' => '1211440166',
                            'total_items' => '410956',
                            'cmd_set' => '588167',
                            'bytes_written' => '35715151',
                            'evictions' => '0',
                            'curr_items' => '31',
                            'pid' => '23599',
                            'limit_maxbytes' => '67108864',
                            'uptime' => '1254',
                            'rusage_user' => '9.857805',
                            'cmd_get' => '838451',
                            'rusage_system' => '34.096988',
                            'version' => '1.2.5',
                            'get_hits' => '581511',
                            'bytes_read' => '46665716',
                            'threads' => '1',
                            'total_connections' => '3104',
                            'get_misses' => '256940'
                          },
                      'sizes' => {
                             '128' => '16',
                             '64' => '15'
                           }
                    }
         },
    'self' => {},
    'total' => {
           'cmd_get' => 838451,
           'bytes' => 2421,
           'get_hits' => 581511,
           'connection_structures' => 56,
           'bytes_read' => 46665716,
           'total_items' => 410956,
           'total_connections' => 3104,
           'cmd_set' => 588167,
           'bytes_written' => 35715151,
           'curr_items' => 31,
           'get_misses' => 256940
         }
       };
```

The statistics are divided up into a number of distinct sections, and then can be requested by adding the type to the stats command. Each statistics output is covered in more detail in the following sections.

- General statistics, see Section 17.2.4.1, “memcached General Statistics”.
- Slab statistics (slabs), see Section 17.2.4.2, “memcached Slabs Statistics”.
- Item statistics (items), see Section 17.2.4.3, “memcached Item Statistics”.
- Size statistics (sizes), see Section 17.2.4.4, “memcached Size Statistics”.
- Detailed status (detail), see Section 17.2.4.5, “memcached Detail Statistics”.

17.2.4.1 memcached General Statistics

The output of the general statistics provides an overview of the performance and use of the memcached instance. The statistics returned by the command and their meaning is shown in the following table.

The following terms are used to define the value type for each statistics value:

- 32u: 32-bit unsigned integer
- 64u: 64-bit unsigned integer
- 32u:32u: Two 32-bit unsigned integers separated by a colon
- String: Character string

Statistic	Data type	Description	Version
pid	32u	Process ID of the memcached instance.	
uptime	32u	Uptime (in seconds) for this memcached instance.	
time	32u	Current time (as epoch).	
version	string	Version string of this instance.	
pointer_size	string	Size of pointers for this host specified in bits (32 or 64).	
rusage_user	32u:32u	Total user time for this instance (seconds:microseconds).	
rusage_system	32u:32u	Total system time for this instance (seconds:microseconds).	
curr_items	32u	Current number of items stored by this instance.	
total_items	32u	Total number of items stored during the life of this instance.	
bytes	64u	Current number of bytes used by this server to store items.	
curr_connections	32u	Current number of open connections.	
total_connections	32u	Total number of connections opened since the server started running.	
connection_structures	32u	Number of connection structures allocated by the server.	
cmd_get	64u	Total number of retrieval requests (get operations).	
cmd_set	64u	Total number of storage requests (set operations).	
get_hits	64u	Number of keys that have been requested and found present.	
get_misses	64u	Number of items that have been requested and not found.	
delete_hits	64u	Number of keys that have been deleted and found present.	1.3.x
delete_misses	64u	Number of items that have been delete and not found.	1.3.x
incr_hits	64u	Number of keys that have been incremented and found present.	1.3.x
incr_misses	64u	Number of items that have been incremented and not found.	1.3.x
decr_hits	64u	Number of keys that have been decremented and found present.	1.3.x

Statistic	Data type	Description	Version
decr_misses	64u	Number of items that have been decremented and not found.	1.3.x
cas_hits	64u	Number of keys that have been compared and swapped and found present.	1.3.x
cas_misses	64u	Number of items that have been compared and swapped and not found.	1.3.x
cas_badvalue	64u	Number of keys that have been compared and swapped, but the comparison (original) value did not match the supplied value.	1.3.x
evictions	64u	Number of valid items removed from cache to free memory for new items.	
bytes_read	64u	Total number of bytes read by this server from network.	
bytes_written	64u	Total number of bytes sent by this server to network.	
limit_maxbytes	32u	Number of bytes this server is permitted to use for storage.	
threads	32u	Number of worker threads requested.	
conn_yields	64u	Number of yields for connections (related to the -R option).	1.4.0

The most useful statistics from those given here are the number of cache hits, misses, and evictions.

A large number of get_misses may just be an indication that the cache is still being populated with information. The number should, over time, decrease in comparison to the number of cache get_hits. If, however, you have a large number of cache misses compared to cache hits after an extended period of execution, it may be an indication that the size of the cache is too small and you either need to increase the total memory size, or increase the number of the memcached instances to improve the hit ratio.

A large number of evictions from the cache, particularly in comparison to the number of items stored is a sign that your cache is too small to hold the amount of information that you regularly want to keep cached. Instead of items being retained in the cache, items are being evicted to make way for new items keeping the turnover of items in the cache high, reducing the efficiency of the cache.

17.2.4.2 memcached Slabs Statistics

To get the slabs statistics, use the stats slabs command, or the API equivalent.

The slab statistics provide you with information about the slabs that have created and allocated for storing information within the cache. You get information both on each individual slab-class and total statistics for the whole slab.

```
STAT 1:chunk_size 104
STAT 1:chunks_per_page 10082
STAT 1:total_pages 1
STAT 1:total_chunks 10082
STAT 1:used_chunks 10081
STAT 1:free_chunks 1
STAT 1:free_chunks_end 10079
STAT 9:chunk_size 696
STAT 9:chunks_per_page 1506
STAT 9:total_pages 63
STAT 9:total_chunks 94878
STAT 9:used_chunks 94878
STAT 9:free_chunks 0
STAT 9:free_chunks_end 0
STAT active_slabs 2
```

```
STAT total_malloced 67083616
END
```

Individual stats for each slab class are prefixed with the slab ID. A unique ID is given to each allocated slab from the smallest size up to the largest. The prefix number indicates the slab class number in relation to the calculated chunk from the specified growth factor. Hence in the example, 1 is the first chunk size and 9 is the 9th chunk allocated size.

The parameters returned for each chunk size and a description of each parameter are provided in the following table.

Statistic	Description	Version
chunk_size	Space allocated to each chunk within this slab class.	
chunks_per_page	Number of chunks within a single page for this slab class.	
total_pages	Number of pages allocated to this slab class.	
total_chunks	Number of chunks allocated to the slab class.	
used_chunks	Number of chunks allocated to an item..	
free_chunks	Number of chunks not yet allocated to items.	
free_chunks_end	Number of free chunks at the end of the last allocated page.	
get_hits	Number of get hits to this chunk	1.3.x
cmd_set	Number of set commands on this chunk	1.3.x
delete_hits	Number of delete hits to this chunk	1.3.x
incr_hits	Number of increment hits to this chunk	1.3.x
decr_hits	Number of decrement hits to this chunk	1.3.x
cas_hits	Number of CAS hits to this chunk	1.3.x
cas_badval	Number of CAS hits on this chunk where the existing value did not match	1.3.x
mem_requested	The true amount of memory of memory requested within this chunk	1.4.1

The following additional statistics cover the information for the entire server, rather than on a chunk by chunk basis:

Statistic	Description	Version
active_slabs	Total number of slab classes allocated.	
total_malloced	Total amount of memory allocated to slab pages.	

The key values in the slab statistics are the chunk_size, and the corresponding total_chunks and used_chunks parameters. These given an indication of the size usage of the chunks within the system. Remember that one key/value pair is placed into a chunk of a suitable size.

From these stats, you can get an idea of your size and chunk allocation and distribution. If you store many items with a number of largely different sizes, consider adjusting the chunk size growth factor to increase in larger steps to prevent chunk and memory wastage. A good indication of a bad growth factor is a high number of different slab classes, but with relatively few chunks actually in use within each slab. Increasing the growth factor creates fewer slab classes and therefore makes better use of the allocated pages.

17.2.4.3 memcached Item Statistics

To get the items statistics, use the stats items command, or the API equivalent.

The items statistics give information about the individual items allocated within a given slab class.

```
STAT items:2:number 1
STAT items:2:age 452
STAT items:2:evicted 0
STAT items:2:evicted_nonzero 0
STAT items:2:evicted_time 2
STAT items:2:outofmemory 0
STAT items:2:tailrepairs 0
...
STAT items:27:number 1
STAT items:27:age 452
STAT items:27:evicted 0
STAT items:27:evicted_nonzero 0
STAT items:27:evicted_time 2
STAT items:27:outofmemory 0
STAT items:27:tailrepairs 0
```

The prefix number against each statistics relates to the corresponding chunk size, as returned by the stats slabs statistics. The result is a display of the number of items stored within each chunk within each slab size, and specific statistics about their age, eviction counts, and out of memory counts. A summary of the statistics is given in the following table.

Statistic	Description	
number	The number of items currently stored in this slab class.	
age	The age of the oldest item within the slab class, in seconds.	
evicted	The number of items evicted to make way for new entries.	
evicted_time	The time of the last evicted entry	
evicted_nonzero	The time of the last evicted non-zero entry	1.4.0
outofmemory	The number of items for this slab class that have triggered an out of memory error (only value when the -M command line option is in effect).	
tailrepairs	Number of times the entries for a particular ID need repairing	

Item level statistics can be used to determine how many items are stored within a given slab and their freshness and recycle rate. You can use this to help identify whether there are certain slab classes that are triggering a much larger number of evictions that others.

17.2.4.4 memcached Size Statistics

To get size statistics, use the stats sizes command, or the API equivalent.

The size statistics provide information about the sizes and number of items of each size within the cache. The information is returned as two columns, the first column is the size of the item (rounded up to the nearest 32 byte boundary), and the second column is the count of the number of items of that size within the cache:

```
96 35
128 38
160 807
192 804
224 410
256 222
288 83
320 39
352 53
384 33
416 64
448 51
```

```
480 30
512 54
544 39
576 10065
```

Caution

Running this statistic locks up your cache as each item is read from the cache and its size calculated. On a large cache, this may take some time and prevent any set or get operations until the process completes.

The item size statistics are useful only to determine the sizes of the objects you are storing. Since the actual memory allocation is relevant only in terms of the chunk size and page size, the information is only useful during a careful debugging or diagnostic session.

17.2.4.5 memcached Detail Statistics

For memcached 1.3.x and higher, you can enable and obtain detailed statistics about the get, set, and del operations on theindividual keys stored in the cache, and determine whether the attempts hit (found) a particular key. These operations are only recorded while the detailed stats analysis is turned on.

To enable detailed statistics, you must send the stats detail on command to the memcached server:

```
$ telnet localhost 11211
Trying 127.0.0.1...
Connected to tiger.
Escape character is '^]'.
stats detail on
OK
```

Individual statistics are recorded for every get, set and del operation on a key, including keys that are not currently stored in the server. For example, if an attempt is made to obtain the value of key abckey and it does not exist, the get operating on the specified key are recorded while detailed statistics are in effect, even if the key is not currently stored. The hits, that is, the number of get or del operations for a key that exists in the server are also counted.

To turn detailed statistics off, send the stats detail off command to the memcached server:

```
$ telnet localhost 11211
Trying 127.0.0.1...
Connected to tiger.
Escape character is '^]'.
stats detail off
OK
```

To obtain the detailed statistics recorded during the process, send the stats detail dump command to the memcached server:

```
stats detail dump
PREFIX hykkey get 0 hit 0 set 1 del 0
PREFIX xyzkey get 0 hit 0 set 1 del 0
PREFIX yukkey get 1 hit 0 set 0 del 0
PREFIX abckey get 3 hit 3 set 1 del 0
END
```

You can use the detailed statistics information to determine whether your memcached clients are using a large number of keys that do not exist in the server by comparing the hit and get or del counts. Because the information is recorded by key, you can also determine whether the failures or operations are clustered around specific keys.

17.2.4.6 Using memcached-tool

The memcached-tool, located within the scripts directory within the memcached source directory. The tool provides convenient access to some reports and statistics from any memcached instance.

The basic format of the command is:

```
shell> ./memcached-tool hostname:port [command]
```

The default output produces a list of the slab allocations and usage. For example:

```
shell> memcached-tool localhost:11211 display
  #  Item_Size  Max_age   Pages   Count   Full?  Evicted Evict_Time OOM
  1      80B        93s       1      20      no        0        0    0
  2     104B        93s       1      16      no        0        0    0
  3     136B      1335s       1      28      no        0        0    0
  4     176B      1335s       1      24      no        0        0    0
  5     224B      1335s       1      32      no        0        0    0
  6     280B      1335s       1      34      no        0        0    0
  7     352B      1335s       1      36      no        0        0    0
  8     440B      1335s       1      46      no        0        0    0
  9     552B      1335s       1      58      no        0        0    0
 10     696B      1335s       1      66      no        0        0    0
 11     872B      1335s       1      89      no        0        0    0
 12     1.1K      1335s       1     112      no        0        0    0
 13     1.3K      1335s       1     145      no        0        0    0
 14     1.7K      1335s       1     123      no        0        0    0
 15     2.1K      1335s       1     198      no        0        0    0
 16     2.6K      1335s       1     199      no        0        0    0
 17     3.3K      1335s       1     229      no        0        0    0
 18     4.1K      1335s       1     248     yes       36        2    0
 19     5.2K      1335s       2     328      no        0        0    0
 20     6.4K      1335s       2     316     yes      387        1    0
 21     8.1K      1335s       3     381     yes      492        1    0
 22    10.1K      1335s       3     303     yes      598        2    0
 23    12.6K      1335s       5     405     yes      605        1    0
 24    15.8K      1335s       6     384     yes      766        2    0
 25    19.7K      1335s       7     357     yes      908      170    0
 26    24.6K      1336s       7     287     yes     1012        1    0
 27    30.8K      1336s       7     231     yes     1193      169    0
 28    38.5K      1336s       4     104     yes     1323      169    0
 29    48.1K      1336s       1      21     yes     1287        1    0
 30    60.2K      1336s       1      17     yes     1093      169    0
 31    75.2K      1337s       1      13     yes      713      168    0
 32    94.0K      1337s       1      10     yes      278      168    0
 33   117.5K      1336s       1       3      no        0        0    0
```

This output is the same if you specify the command as display:

```
shell> memcached-tool localhost:11211 display
  #  Item_Size  Max_age   Pages   Count   Full?  Evicted Evict_Time OOM
  1      80B        93s       1      20      no        0        0    0
  2     104B        93s       1      16      no        0        0    0
...
```

The output shows a summarized version of the output from the slabs statistics. The columns provided in the output are shown below:

- #: The slab number
- Item_Size: The size of the slab
- Max_age: The age of the oldest item in the slab
- Pages: The number of pages allocated to the slab

- Count: The number of items in this slab
- Full?: Whether the slab is fully populated
- Evicted: The number of objects evicted from this slab
- Evict_Time: The time (in seconds) since the last eviction
- OOM: The number of items that have triggered an out of memory error

You can also obtain a dump of the general statistics for the server using the stats command:

```
shell> memcached-tool localhost:11211 stats
#localhost:11211   Field       Value
         accepting_conns           1
                   bytes         162
              bytes_read         485
           bytes_written        6820
              cas_badval           0
                cas_hits           0
              cas_misses           0
               cmd_flush           0
                 cmd_get           4
                 cmd_set           2
             conn_yields           0
   connection_structures          11
        curr_connections          10
              curr_items           2
               decr_hits           0
             decr_misses           1
             delete_hits           0
           delete_misses           0
               evictions           0
                get_hits           4
              get_misses           0
               incr_hits           0
             incr_misses           2
          limit_maxbytes    67108864
     listen_disabled_num           0
                     pid       12981
            pointer_size          32
          rusage_system    0.013911
            rusage_user    0.011876
                threads           4
                   time  1255518565
       total_connections          20
             total_items           2
                  uptime         880
                 version       1.4.2
```

17.2.5 memcached FAQ

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17.2.5.1 Can memcached be run on a Windows environment?

No. Currently memcached is available only on the Unix/Linux platform. There is an unofficial port available, see http://www.codeplex.com/memcachedproviders.

17.2.5.2 What is the maximum size of an object you can store in memcached? Is that configurable?

The default maximum object size is 1MB. In memcached 1.4.2 and later, you can change the maximum size of an object using the `-I` command line option.

For versions before this, to increase this size, you have to re-compile memcached. You can modify the value of the `POWER_BLOCK` within the `slabs.c` file within the source.

In memcached 1.4.2 and higher, you can configure the maximum supported object size by using the `-I` command-line option. For example, to increase the maximum object size to 5MB:

```
$ memcached -I 5m
```

If an object is larger than the maximum object size, you must manually split it. memcached is very simple: you give it a key and some data, it tries to cache it in RAM. If you try to store more than the default maximum size, the value is just truncated for speed reasons.

17.2.5.3 Is it true memcached will be much more effective with db-read-intensive applications than with db-write-intensive applications?

Yes. memcached plays no role in database writes, it is a method of caching data already read from the database in RAM.

17.2.5.4 Is there any overhead in not using persistent connections? If persistent is always recommended, what are the downsides (for example, locking up)?

If you don't use persistent connections when communicating with memcached, there will be a small increase in the latency of opening the connection each time. The effect is comparable to use nonpersistent connections with MySQL.

In general, the chance of locking or other issues with persistent connections is minimal, because there is very little locking within memcached. If there is a problem, eventually your request will time out and return no result, so your application will need to load from MySQL again.

17.2.5.5 How is an event such as a crash of one of the memcached servers handled by the memcached client?

There is no automatic handling of this. If your client fails to get a response from a server, code a fallback mechanism to load the data from the MySQL database.

The client APIs all provide the ability to add and remove memcached instances on the fly. If within your application you notice that memcached server is no longer responding, you can remove the server from the list of servers, and keys will automatically be redistributed to another memcached server in the list. If retaining the cache content on all your servers is important, make sure you use an API that supports a consistent hashing algorithm. For more information, see Section 17.2.2.5, “memcached Hashing/Distribution Types”.

17.2.5.6. What is a recommended hardware configuration for a memcached server?

memcached has a very low processing overhead. All that is required is spare physical RAM capacity. A memcached server does not require a dedicated machine. If you have web, application, or database servers that have spare RAM capacity, then use them with memcached.

To build and deploy a dedicated memcached server, use a relatively low-power CPU, lots of RAM, and one or more Gigabit Ethernet interfaces.

17.2.5.7. Is memcached more effective for video and audio as opposed to textual read/writes?

memcached works equally well for all kinds of data. To memcached, any value you store is just a stream of data. Remember, though, that the maximum size of an object you can store in memcached is 1MB, but can be configured to be larger by using the -I option in memcached 1.4.2 and later, or by modifying the source in versions before 1.4.2. If you plan on using memcached with audio and video content, you will probably want to increase the maximum object size. Also remember that memcached is a solution for caching information for reading. It shouldn't be used for writes, except when updating the information in the cache.

17.2.5.8. Can memcached work with ASPX?

There are ports and interfaces for many languages and environments. ASPX relies on an underlying language such as C# or VisualBasic, and if you are using ASP.NET then there is a C# memcached library. For more information, see https://sourceforge.net/projects/memcacheddotnet/.

17.2.5.9. How expensive is it to establish a memcache connection? Should those connections be pooled?

Opening the connection is relatively inexpensive, because there is no security, authentication or other handshake taking place before you can start sending requests and getting results. Most APIs support a persistent connection to a memcached instance to reduce the latency. Connection pooling would depend on the API you are using, but if you are communicating directly over TCP/IP, then connection pooling would provide some small performance benefit.

17.2.5.10. How is the data handled when the memcached server is down?

The behavior is entirely application dependent. Most applications fall back to loading the data from the database (just as if they were updating the memcached information). If you are using multiple memcached servers, you might also remove a downed server from the list to prevent it from affecting performance. Otherwise, the client will still attempt to communicate with the memcached server that corresponds to the key you are trying to load.

17.2.5.11. How are auto-increment columns in the MySQL database coordinated across multiple instances of memcached?

They aren't. There is no relationship between MySQL and memcached unless your application (or, if you are using the MySQL UDFs for memcached, your database definition) creates one.

If you are storing information based on an auto-increment key into multiple instances of memcached, the information is only stored on one of the memcached instances anyway. The client uses the key value to determine which memcached instance to store the information. It doesn't store the same information across all the instances, as that would be a waste of cache memory.

17.2.5.12. Is compression available?

Yes. Most of the client APIs support some sort of compression, and some even allow you to specify the threshold at which a value is deemed appropriate for compression during storage.

17.2.5.13. Can we implement different types of memcached as different nodes in the same server, so can there be deterministic and non-deterministic in the same server?

Yes. You can run multiple instances of memcached on a single server, and in your client configuration you choose the list of servers you want to use.

17.2.5.14 What are best practices for testing an implementation, to ensure that it improves performance, and to measure the impact of memcached configuration changes? And would you recommend keeping the configuration very simple to start?

The best way to test the performance is to start up a memcached instance. First, modify your application so that it stores the data just before the data is about to be used or displayed into memcached. Since the APIs handle the serialization of the data, it should just be a one-line modification to your code. Then, modify the start of the process that would normally load that information from MySQL with the code that requests the data from memcached. If the data cannot be loaded from memcached, default to the MySQL process.

All of the changes required will probably amount to just a few lines of code. To get the best benefit, make sure you cache entire objects (for example, all the components of a web page, blog post, discussion thread, and so on), rather than using memcached as a simple cache of individual rows of MySQL tables.

Keeping the configuration simple at the start, or even over the long term, is easy with memcached. Once you have the basic structure up and running, often the only ongoing change is to add more servers into the list of servers used by your applications. You don't need to manage the memcached servers, and there is no complex configuration; just add more servers to the list and let the client API and the memcached servers make the decisions.

Chapter 18 Replication

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Replication enables data from one MySQL database server (the master) to be copied to one or more MySQL database servers (the slaves). Replication is asynchronous by default; slaves do not need to be connected permanently to receive updates from the master. Depending on the configuration, you can replicate all databases, selected databases, or even selected tables within a database.

Advantages of replication in MySQL include:

- Scale-out solutions - spreading the load among multiple slaves to improve performance. In this environment, all writes and updates must take place on the master server. Reads, however, may take place on one or more slaves. This model can improve the performance of writes (since the master is dedicated to updates), while dramatically increasing read speed across an increasing number of slaves.

- Data security - because data is replicated to the slave, and the slave can pause the replication process, it is possible to run backup services on the slave without corrupting the corresponding master data.

- Analytics - live data can be created on the master, while the analysis of the information can take place on the slave without affecting the performance of the master.

- Long-distance data distribution - you can use replication to create a local copy of data for a remote site to use, without permanent access to the master.

For information on how to use replication in such scenarios, see Section 18.3, “Replication Solutions”.

MySQL 8.0 supports different methods of replication. The traditional method is based on replicating events from the master's binary log, and requires the log files and positions in them to be synchronized between master and slave. The newer method based on *global transaction identifiers* (GTIDs) is transactional and therefore does not require working with log files or positions within these files, which greatly simplifies many common replication tasks. Replication using GTIDs guarantees consistency between master and slave as long as all transactions committed on the master have also been applied on the slave. For more information about GTIDs and GTID-based replication in MySQL, see Section 18.1.3, “Replication with Global Transaction Identifiers”. For information on using binary log file position based replication, see Section 18.1, “Configuring Replication”.

Replication in MySQL supports different types of synchronization. The original type of synchronization is one-way, asynchronous replication, in which one server acts as the master, while one or more other servers act as slaves. In MySQL 8.0, semisynchronous replication is supported in addition to the built-in asynchronous replication. With semisynchronous replication, a commit performed on the master blocks before returning to the session that performed the transaction until at least one slave acknowledges that it has received and logged the events for the transaction; see Section 18.3.10, “Semisynchronous Replication”. MySQL 8.0 also supports delayed replication such that a slave server deliberately lags behind the master by at least a specified amount of time; see Section 18.3.11, “Delayed Replication”. For scenarios where *synchronous* replication is required, use MySQL Cluster (see MySQL NDB Cluster 7.5 and NDB Cluster 7.6).

There are a number of solutions available for setting up replication between servers, and the best method to use depends on the presence of data and the engine types you are using. For more information on the available options, see Section 18.1.2, “Setting Up Binary Log File Position Based Replication”.

There are two core types of replication format, Statement Based Replication (SBR), which replicates entire SQL statements, and Row Based Replication (RBR), which replicates only the changed rows. You can also use a third variety, Mixed Based Replication (MBR). For more information on the different replication formats, see Section 18.2.1, “Replication Formats”.

Replication is controlled through a number of different options and variables. For more information, see Section 18.1.6, “Replication and Binary Logging Options and Variables”.

You can use replication to solve a number of different problems, including performance, supporting the backup of different databases, and as part of a larger solution to alleviate system failures. For information on how to address these issues, see Section 18.3, “Replication Solutions”.

For notes and tips on how different data types and statements are treated during replication, including details of replication features, version compatibility, upgrades, and potential problems and their resolution, see Section 18.4, “Replication Notes and Tips”. For answers to some questions often asked by those who are new to MySQL Replication, see Section A.13, “MySQL 8.0 FAQ: Replication”.

For detailed information on the implementation of replication, how replication works, the process and contents of the binary log, background threads and the rules used to decide how statements are recorded and replicated, see Section 18.2, “Replication Implementation”.

18.1 Configuring Replication

This section describes how to configure the different types of replication available in MySQL and includes the setup and configuration required for a replication environment, including step-by-step instructions for creating a new replication environment. The major components of this section are:

- For a guide to setting up two or more servers for replication using binary log file positions, Section 18.1.2, “Setting Up Binary Log File Position Based Replication”, deals with the configuration of the servers and provides methods for copying data between the master and slaves.

- For a guide to setting up two or more servers for replication using GTID transactions, Section 18.1.3, “Replication with Global Transaction Identifiers”, deals with the configuration of the servers.
- Events in the binary log are recorded using a number of formats. These are referred to as statement-based replication (SBR) or row-based replication (RBR). A third type, mixed-format replication (MIXED), uses SBR or RBR replication automatically to take advantage of the benefits of both SBR and RBR formats when appropriate. The different formats are discussed in Section 18.2.1, “Replication Formats”.
- Detailed information on the different configuration options and variables that apply to replication is provided in Section 18.1.6, “Replication and Binary Logging Options and Variables”.
- Once started, the replication process should require little administration or monitoring. However, for advice on common tasks that you may want to execute, see Section 18.1.7, “Common Replication Administration Tasks”.

18.1.1 Binary Log File Position Based Replication Configuration Overview

This section describes replication between MySQL servers based on the binary log file position method, where the MySQL instance operating as the master (the source of the database changes) writes updates and changes as “events” to the binary log. The information in the binary log is stored in different logging formats according to the database changes being recorded. Slaves are configured to read the binary log from the master and to execute the events in the binary log on the slave's local database.

Each slave receives a copy of the entire contents of the binary log. It is the responsibility of the slave to decide which statements in the binary log should be executed. Unless you specify otherwise, all events in the master binary log are executed on the slave. If required, you can configure the slave to process only events that apply to particular databases or tables.

Important

You cannot configure the master to log only certain events.

Each slave keeps a record of the binary log coordinates: the file name and position within the file that it has read and processed from the master. This means that multiple slaves can be connected to the master and executing different parts of the same binary log. Because the slaves control this process, individual slaves can be connected and disconnected from the server without affecting the master's operation. Also, because each slave records the current position within the binary log, it is possible for slaves to be disconnected, reconnect and then resume processing.

The master and each slave must be configured with a unique ID (using the `server-id` option). In addition, each slave must be configured with information about the master host name, log file name, and position within that file. These details can be controlled from within a MySQL session using the `CHANGE MASTER TO` statement on the slave. The details are stored within the slave's master info repository, which can be either a file or a table (see Section 18.2.4, “Replication Relay and Status Logs”).

18.1.2 Setting Up Binary Log File Position Based Replication

This section describes how to set up a MySQL server to use binary log file position based replication. There are a number of different methods for setting up replication, and the exact method to use depends on how you are setting up replication, and whether you already have data within your master database.

There are some generic tasks that are common to all setups:

- On the master, you must enable binary logging and configure a unique server ID. This might require a server restart. See Section 18.1.2.1, “Setting the Replication Master Configuration”.

- On each slave that you want to connect to the master, you must configure a unique server ID. This might require a server restart. See Section 18.1.2.2, “Setting the Replication Slave Configuration”.
- Optionally, create a separate user for your slaves to use during authentication with the master when reading the binary log for replication. See Section 18.1.2.3, “Creating a User for Replication”.
- Before creating a data snapshot or starting the replication process, on the master you should record the current position in the binary log. You need this information when configuring the slave so that the slave knows where within the binary log to start executing events. See Section 18.1.2.4, “Obtaining the Replication Master Binary Log Coordinates”.
- If you already have data on the master and want to use it to synchronize the slave, you need to create a data snapshot to copy the data to the slave. The storage engine you are using has an impact on how you create the snapshot. When you are using `MyISAM`, you must stop processing statements on the master to obtain a read-lock, then obtain its current binary log coordinates and dump its data, before permitting the master to continue executing statements. If you do not stop the execution of statements, the data dump and the master status information will not match, resulting in inconsistent or corrupted databases on the slaves. For more information on replicating a `MyISAM` master, see Section 18.1.2.4, “Obtaining the Replication Master Binary Log Coordinates”. If you are using `InnoDB`, you do not need a read-lock and a transaction that is long enough to transfer the data snapshot is sufficient. For more information, see Section 15.18, “InnoDB and MySQL Replication”.
- Configure the slave with settings for connecting to the master, such as the host name, login credentials, and binary log file name and position. See Section 18.1.2.7, “Setting the Master Configuration on the Slave”.

Note

Certain steps within the setup process require the `SUPER` privilege. If you do not have this privilege, it might not be possible to enable replication.

After configuring the basic options, select your scenario:

- To set up replication for a fresh installation of a master and slaves that contain no data, see Setting Up Replication with New Master and Slaves.
- To set up replication of a new master using the data from an existing MySQL server, see Setting Up Replication with Existing Data.
- To add replication slaves to an existing replication environment, see Section 18.1.2.8, “Adding Slaves to a Replication Environment”.

Before administering MySQL replication servers, read this entire chapter and try all statements mentioned in Section 13.4.1, “SQL Statements for Controlling Master Servers”, and Section 13.4.2, “SQL Statements for Controlling Slave Servers”. Also familiarize yourself with the replication startup options described in Section 18.1.6, “Replication and Binary Logging Options and Variables”.

18.1.2.1 Setting the Replication Master Configuration

To configure a master to use binary log file position based replication, you must enable binary logging and establish a unique server ID. If this has not already been done, a server restart is required.

Binary logging *must* be enabled on the master because the binary log is the basis for replicating changes from the master to its slaves. If binary logging is not enabled on the master using the `log-bin` option, replication is not possible.

Each server within a replication group must be configured with a unique server ID. This ID is used to identify individual servers within the group, and must be a positive integer between 1 and (2^{32})−1. How you organize and select the numbers is your choice.

To configure the binary log and server ID options, shut down the MySQL server and edit the `my.cnf` or `my.ini` file. Within the `[mysqld]` section of the configuration file, add the `log-bin` and `server-id`

options. If these options already exist, but are commented out, uncomment the options and alter them according to your needs. For example, to enable binary logging using a log file name prefix of `mysql-bin`, and configure a server ID of 1, use these lines:

```
[mysqld]
log-bin=mysql-bin
server-id=1
```

After making the changes, restart the server.

Note

The following options have an impact on this procedure:

- if you omit `server-id` (or set it explicitly to its default value of 0), the master refuses any connections from slaves.
- For the greatest possible durability and consistency in a replication setup using `InnoDB` with transactions, you should use `innodb_flush_log_at_trx_commit=1` and `sync_binlog=1` in the master `my.cnf` file.
- Ensure that the `skip-networking` option is not enabled on your replication master. If networking has been disabled, the slave can not communicate with the master and replication fails.

18.1.2.2 Setting the Replication Slave Configuration

Each replication slave *must* have a unique server ID. If this has not already been done, this part of slave setup requires a server restart.

If the slave server ID is not already set, or the current value conflicts with the value that you have chosen for the master server, shut down the slave server and edit the `[mysqld]` section of the configuration file to specify a unique server ID. For example:

```
[mysqld]
server-id=2
```

After making the changes, restart the server.

If you are setting up multiple slaves, each one must have a unique `server-id` value that differs from that of the master and from any of the other slaves.

Note

If you omit `server-id` (or set it explicitly to its default value of 0), the slave refuses to connect to a master.

You do not have to enable binary logging on the slave for replication to be set up. However, if you enable binary logging on the slave, you can use the slave's binary log for data backups and crash recovery, and also use the slave as part of a more complex replication topology. For example, where this slave then acts as a master to other slaves.

18.1.2.3 Creating a User for Replication

Each slave connects to the master using a MySQL user name and password, so there must be a user account on the master that the slave can use to connect. Any account can be used for this operation, providing it has been granted the `REPLICATION SLAVE` privilege. You can choose to create a different account for each slave, or connect to the master using the same account for each slave.

Although you do not have to create an account specifically for replication, you should be aware that the replication user name and password are stored in plain text in the master info repository file or table

(see Section 18.2.4.2, “Slave Status Logs”). Therefore, you may want to create a separate account that has privileges only for the replication process, to minimize the possibility of compromise to other accounts.

To create a new account, use `CREATE USER`. To grant this account the privileges required for replication, use the `GRANT` statement. If you create an account solely for the purposes of replication, that account needs only the `REPLICATION SLAVE` privilege. For example, to set up a new user, `repl`, that can connect for replication from any host within the `mydomain.com` domain, issue these statements on the master:

```
mysql> CREATE USER 'repl'@'%.mydomain.com' IDENTIFIED BY 'slavepass';
mysql> GRANT REPLICATION SLAVE ON *.* TO 'repl'@'%.mydomain.com';
```

See Section 13.7.1, “Account Management Statements”, for more information on statements for manipulation of user accounts.

18.1.2.4 Obtaining the Replication Master Binary Log Coordinates

To configure the slave to start the replication process at the correct point, you need to note the master's current coordinates within its binary log.

Warning

This procedure uses `FLUSH TABLES WITH READ LOCK`, which blocks `COMMIT` operations for `InnoDB` tables.

To obtain the master binary log coordinates, follow these steps:

1. Start a session on the master by connecting to it with the command-line client, and flush all tables and block write statements by executing the `FLUSH TABLES WITH READ LOCK` statement:

   ```
   mysql> FLUSH TABLES WITH READ LOCK;
   ```

   Warning

   Leave the client from which you issued the `FLUSH TABLES` statement running so that the read lock remains in effect. If you exit the client, the lock is released.

2. In a different session on the master, use the `SHOW MASTER STATUS` statement to determine the current binary log file name and position:

   ```
   mysql > SHOW MASTER STATUS;
   +------------------+----------+--------------+------------------+
   | File             | Position | Binlog_Do_DB | Binlog_Ignore_DB |
   +------------------+----------+--------------+------------------+
   | mysql-bin.000003 | 73       | test         | manual,mysql     |
   +------------------+----------+--------------+------------------+
   ```

   The `File` column shows the name of the log file and the `Position` column shows the position within the file. In this example, the binary log file is `mysql-bin.000003` and the position is 73. Record these values. You need them later when you are setting up the slave. They represent the replication coordinates at which the slave should begin processing new updates from the master.

   If the master has been running previously without binary logging enabled, the log file name and position values displayed by `SHOW MASTER STATUS` or `mysqldump --master-data` will be empty. In that case, the values that you need to use later when specifying the slave's log file and position are the empty string (`''`) and `4`.

You now have the information you need to enable the slave to start reading from the binary log in the correct place to start replication.

The next step depends on whether you have existing data on the master. Choose one of the following options:

- If you have existing data that needs be to synchronized with the slave before you start replication, leave the client running so that the lock remains in place. This prevents any further changes being made, so that the data copied to the slave is in synchrony with the master. Proceed to Section 18.1.2.5, “Choosing a Method for Data Snapshots”.
- If you are setting up a new master and slave replication group, you can exit the first session to release the read lock. See Setting Up Replication with New Master and Slaves for how to proceed.

18.1.2.5 Choosing a Method for Data Snapshots

If the master database contains existing data it is necessary to copy this data to each slave. There are different ways to dump the data from the master database. The following sections describe possible options.

To select the appropriate method of dumping the database, choose between these options:

- Use the `mysqldump` tool to create a dump of all the databases you want to replicate. This is the recommended method, especially when using `InnoDB`.
- If your database is stored in binary portable files, you can copy the raw data files to a slave. This can be more efficient than using `mysqldump` and importing the file on each slave, because it skips the overhead of updating indexes as the `INSERT` statements are replayed. With storage engines such as `InnoDB` this is not recommended.

Creating a Data Snapshot Using mysqldump

To create a snapshot of the data in an existing master database, use the `mysqldump` tool. Once the data dump has been completed, import this data into the slave before starting the replication process.

The following example dumps all databases to a file named `dbdump.db`, and includes the `--master-data` option which automatically appends the `CHANGE MASTER TO` statement required on the slave to start the replication process:

```
shell> mysqldump --all-databases --master-data > dbdump.db
```

Note

If you do not use `--master-data`, then it is necessary to lock all tables in a separate session manually. See Section 18.1.2.4, “Obtaining the Replication Master Binary Log Coordinates”.

It is possible to exclude certain databases from the dump using the `mysqldump` tool. If you want to choose which databases to include in the dump, do not use `--all-databases`. Choose one of these options:

- Exclude all the tables in the database using `--ignore-table` option.
- Name only those databases which you want dumped using the `--databases` option.

For more information, see Section 4.5.4, “`mysqldump` — A Database Backup Program”.

To import the data, either copy the dump file to the slave, or access the file from the master when connecting remotely to the slave.

Creating a Data Snapshot Using Raw Data Files

This section describes how to create a data snapshot using the raw files which make up the database. Employing this method with a table using a storage engine that has complex caching or logging algorithms requires extra steps to produce a perfect “point in time” snapshot: the initial copy command

could leave out cache information and logging updates, even if you have acquired a global read lock. How the storage engine responds to this depends on its crash recovery abilities.

If you use `InnoDB` tables, you can use the `mysqlbackup` command from the MySQL Enterprise Backup component to produce a consistent snapshot. This command records the log name and offset corresponding to the snapshot to be used on the slave. MySQL Enterprise Backup is a commercial product that is included as part of a MySQL Enterprise subscription. See Section 29.2, "MySQL Enterprise Backup Overview" for detailed information.

This method also does not work reliably if the master and slave have different values for `ft_stopword_file`, `ft_min_word_len`, or `ft_max_word_len` and you are copying tables having full-text indexes.

Assuming the above exceptions do not apply to your database, use the cold backup technique to obtain a reliable binary snapshot of `InnoDB` tables: do a slow shutdown of the MySQL Server, then copy the data files manually.

To create a raw data snapshot of `MyISAM` tables when your MySQL data files exist on a single file system, you can use standard file copy tools such as `cp` or `copy`, a remote copy tool such as `scp` or `rsync`, an archiving tool such as `zip` or `tar`, or a file system snapshot tool such as `dump`. If you are replicating only certain databases, copy only those files that relate to those tables. For `InnoDB`, all tables in all databases are stored in the system tablespace files, unless you have the `innodb_file_per_table` option enabled.

The following files are not required for replication:

- Files relating to the `mysql` database.
- The master info repository file, if used (see Section 18.2.4, "Replication Relay and Status Logs").
- The master's binary log files.
- Any relay log files.

Depending on whether you are using `InnoDB` tables or not, choose one of the following:

If you are using `InnoDB` tables, and also to get the most consistent results with a raw data snapshot, shut down the master server during the process, as follows:

1. Acquire a read lock and get the master's status. See Section 18.1.2.4, "Obtaining the Replication Master Binary Log Coordinates".
2. In a separate session, shut down the master server:

```
shell> mysqladmin shutdown
```

3. Make a copy of the MySQL data files. The following examples show common ways to do this. You need to choose only one of them:

```
shell> tar cf /tmp/db.tar ./data
shell> zip -r /tmp/db.zip ./data
shell> rsync --recursive ./data /tmp/dbdata
```

4. Restart the master server.

If you are not using `InnoDB` tables, you can get a snapshot of the system from a master without shutting down the server as described in the following steps:

1. Acquire a read lock and get the master's status. See Section 18.1.2.4, "Obtaining the Replication Master Binary Log Coordinates".
2. Make a copy of the MySQL data files. The following examples show common ways to do this. You need to choose only one of them:

```
shell> tar cf /tmp/db.tar ./data
shell> zip -r /tmp/db.zip ./data
shell> rsync --recursive ./data /tmp/dbdata
```

3. In the client where you acquired the read lock, release the lock:

```
mysql> UNLOCK TABLES;
```

Once you have created the archive or copy of the database, copy the files to each slave before starting the slave replication process.

18.1.2.6 Setting Up Replication Slaves

The following sections describe how to set up slaves. Before you proceed, ensure that you have:

- Configured the MySQL master with the necessary configuration properties. See Section 18.1.2.1, “Setting the Replication Master Configuration”.
- Obtained the master status information. See Section 18.1.2.4, “Obtaining the Replication Master Binary Log Coordinates”.
- On the master, released the read lock:

```
mysql> UNLOCK TABLES;
```

- On the slave, edited the MySQL configuration. See Section 18.1.2.2, “Setting the Replication Slave Configuration”.

The next steps depend on whether you have existing data to import to the slave or not. See Section 18.1.2.5, “Choosing a Method for Data Snapshots” for more information. Choose one of the following:

- If you do not have a snapshot of a database to import, see Setting Up Replication with New Master and Slaves.
- If you have a snapshot of a database to import, see Setting Up Replication with Existing Data.

Setting Up Replication with New Master and Slaves

When there is no snapshot of a previous database to import, configure the slave to start the replication from the new master.

To set up replication between a master and a new slave:

1. Start up the MySQL slave.
2. Execute a `CHANGE MASTER TO` statement to set the master replication server configuration. See Section 18.1.2.7, “Setting the Master Configuration on the Slave”.

Perform these slave setup steps on each slave.

This method can also be used if you are setting up new servers but have an existing dump of the databases from a different server that you want to load into your replication configuration. By loading the data into a new master, the data is automatically replicated to the slaves.

If you are setting up a new replication environment using the data from a different existing database server to create a new master, run the dump file generated from that server on the new master. The database updates are automatically propagated to the slaves:

```
shell> mysql -h master < fulldb.dump
```

Setting Up Replication with Existing Data

When setting up replication with existing data, transfer the snapshot from the master to the slave before starting replication. The process for importing data to the slave depends on how you created the snapshot of data on the master.

Choose one of the following:

If you used `mysqldump`:

1. Start the slave, using the `--skip-slave-start` option so that replication does not start.
2. Import the dump file:

```
shell> mysql < fulldb.dump
```

If you created a snapshot using the raw data files:

1. Extract the data files into your slave data directory. For example:

```
shell> tar xvf dbdump.tar
```

   You may need to set permissions and ownership on the files so that the slave server can access and modify them.

2. Start the slave, using the `--skip-slave-start` option so that replication does not start.
3. Configure the slave with the replication coordinates from the master. This tells the slave the binary log file and position within the file where replication needs to start. Also, configure the slave with the login credentials and host name of the master. For more information on the `CHANGE MASTER TO` statement required, see Section 18.1.2.7, “Setting the Master Configuration on the Slave”.
4. Start the slave threads:

```
mysql> START SLAVE;
```

After you have performed this procedure, the slave connects to the master and replicates any updates that have occurred on the master since the snapshot was taken.

If you have forgotten to set the `server-id` option for the master, slaves cannot connect to it.

If you have forgotten to set the `server-id` option for the slave, you get the following error in the slave's error log:

```
Warning: You should set server-id to a non-0 value if master_host
is set; we will force server id to 2, but this MySQL server will
not act as a slave.
```

You also find error messages in the slave's error log if it is not able to replicate for any other reason.

The slave uses information stored in its master info repository to keep track of how much of the master's binary log it has processed. The repository can be in the form of files or a table, as determined by the value set for `--master-info-repository`. When a slave runs with `--master-info-repository=FILE`, you can find in its data directory two files, named `master.info` and `relay-log.info`. If `--master-info-repository=TABLE` instead, this information is saved in the table `master_slave_info` in the `mysql` database. In either case, do *not* remove or edit the files or table unless you know exactly what you are doing and fully understand the implications. Even in that case, it is preferred that you use the `CHANGE MASTER TO` statement to change replication parameters. The slave can use the values specified in the statement to update the status files automatically. See Section 18.2.4, “Replication Relay and Status Logs”, for more information.

Note

The contents of the master info repository override some of the server options specified on the command line or in `my.cnf`. See Section 18.1.6, “Replication and Binary Logging Options and Variables”, for more details.

A single snapshot of the master suffices for multiple slaves. To set up additional slaves, use the same master snapshot and follow the slave portion of the procedure just described.

18.1.2.7 Setting the Master Configuration on the Slave

To set up the slave to communicate with the master for replication, configure the slave with the necessary connection information. To do this, execute the following statement on the slave, replacing the option values with the actual values relevant to your system:

```
mysql> CHANGE MASTER TO
    ->     MASTER_HOST='master_host_name',
    ->     MASTER_USER='replication_user_name',
    ->     MASTER_PASSWORD='replication_password',
    ->     MASTER_LOG_FILE='recorded_log_file_name',
    ->     MASTER_LOG_POS=recorded_log_position;
```

Note

Replication cannot use Unix socket files. You must be able to connect to the master MySQL server using TCP/IP.

The `CHANGE MASTER TO` statement has other options as well. For example, it is possible to set up secure replication using SSL. For a full list of options, and information about the maximum permissible length for the string-valued options, see Section 13.4.2.1, “CHANGE MASTER TO Syntax”.

18.1.2.8 Adding Slaves to a Replication Environment

You can add another slave to an existing replication configuration without stopping the master. Instead, set up the new slave by making a copy of an existing slave, except that you configure the new slave with a different `server-id` value.

To duplicate an existing slave:

1. Shut down the existing slave:

   ```
   shell> mysqladmin shutdown
   ```

2. Copy the data directory from the existing slave to the new slave. You can do this by creating an archive using `tar` or `WinZip`, or by performing a direct copy using a tool such as `cp` or `rsync`. Ensure that you also copy the log files and relay log files.

   A common problem that is encountered when adding new replication slaves is that the new slave fails with a series of warning and error messages like these:

   ```
   071118 16:44:10 [Warning] Neither --relay-log nor --relay-log-index were used; so
   replication may break when this MySQL server acts as a slave and has his hostname
   changed!! Please use '--relay-log=new_slave_hostname-relay-bin' to avoid this problem.
   071118 16:44:10 [ERROR] Failed to open the relay log './old_slave_hostname-relay-bin.003525'
   (relay_log_pos 22940879)
   071118 16:44:10 [ERROR] Could not find target log during relay log initialization
   071118 16:44:10 [ERROR] Failed to initialize the master info structure
   ```

   This situation can occur if the `--relay-log` option is not specified, as the relay log files contain the host name as part of their file names. This is also true of the relay log index file if the `--relay-log-index` option is not used. See Section 18.1.6, “Replication and Binary Logging Options and Variables”, for more information about these options.

To avoid this problem, use the same value for `--relay-log` on the new slave that was used on the existing slave. If this option was not set explicitly on the existing slave, use *`existing_slave_hostname`*`-relay-bin`. If this is not possible, copy the existing slave's relay log index file to the new slave and set the `--relay-log-index` option on the new slave to match what was used on the existing slave. If this option was not set explicitly on the existing slave, use *`existing_slave_hostname`*`-relay-bin.index`. Alternatively, if you have already tried to start the new slave after following the remaining steps in this section and have encountered errors like those described previously, then perform the following steps:

a. If you have not already done so, issue a `STOP SLAVE` on the new slave.

   If you have already started the existing slave again, issue a `STOP SLAVE` on the existing slave as well.

b. Copy the contents of the existing slave's relay log index file into the new slave's relay log index file, making sure to overwrite any content already in the file.

c. Proceed with the remaining steps in this section.

3. Copy the master info and relay log info repositories (see Section 18.2.4, “Replication Relay and Status Logs”) from the existing slave to the new slave. These hold the current log coordinates for the master's binary log and the slave's relay log.

4. Start the existing slave.

5. On the new slave, edit the configuration and give the new slave a unique `server-id` not used by the master or any of the existing slaves.

6. Start the new slave. The slave uses the information in its master info repository to start the replication process.

18.1.3 Replication with Global Transaction Identifiers

This section explains transaction-based replication using *global transaction identifiers* (GTIDs). When using GTIDs, each transaction can be identified and tracked as it is committed on the originating server and applied by any slaves; this means that it is not necessary when using GTIDs to refer to log files or positions within those files when starting a new slave or failing over to a new master, which greatly simplifies these tasks. Because GTID-based replication is completely transaction-based, it is simple to determine whether masters and slaves are consistent; as long as all transactions committed on a master are also committed on a slave, consistency between the two is guaranteed. You can use either statement-based or row-based replication with GTIDs (see Section 18.2.1, “Replication Formats”); however, for best results, we recommend that you use the row-based format.

This section discusses the following topics:

- How GTIDs are defined and created, and how they are represented in the MySQL Server (see Section 18.1.3.1, “GTID Concepts”).
- A general procedure for setting up and starting GTID-based replication (see Section 18.1.3.2, “Setting Up Replication Using GTIDs”).
- Suggested methods for provisioning new replication servers when using GTIDs (see Section 18.1.3.3, “Using GTIDs for Failover and Scaleout”).
- Restrictions and limitations that you should be aware of when using GTID-based replication (see Section 18.1.3.4, “Restrictions on Replication with GTIDs”).

For information about MySQL Server options and variables relating to GTID-based replication, see Section 18.1.6.5, “Global Transaction ID Options and Variables”. See also Section 12.17, “Functions Used with Global Transaction IDs”, which describes SQL functions supported by MySQL 8.0 for use with GTIDs.

18.1.3.1 GTID Concepts

A global transaction identifier (GTID) is a unique identifier created and associated with each transaction committed on the server of origin (master). This identifier is unique not only to the server on which it originated, but is unique across all servers in a given replication setup. There is a 1-to-1 mapping between all transactions and all GTIDs.

The following paragraphs provide a basic description of GTIDs. More advanced concepts are covered later in the following sections:

- GTID Sets
- mysql.gtid_executed Table
- mysql.gtid_executed Table Compression

A GTID is represented as a pair of coordinates, separated by a colon character (`:`), as shown here:

```
GTID = source_id:transaction_id
```

The *`source_id`* identifies the originating server. Normally, the server's `server_uuid` is used for this purpose. The *`transaction_id`* is a sequence number determined by the order in which the transaction was committed on this server; for example, the first transaction to be committed has `1` as its *`transaction_id`*, and the tenth transaction to be committed on the same originating server is assigned a *`transaction_id`* of `10`. It is not possible for a transaction to have `0` as a sequence number in a GTID. For example, the twenty-third transaction to be committed originally on the server with the UUID `3E11FA47-71CA-11E1-9E33-C80AA9429562` has this GTID:

```
3E11FA47-71CA-11E1-9E33-C80AA9429562:23
```

This format is used to represent GTIDs in the output of statements such as `SHOW SLAVE STATUS` as well as in the binary log. They can also be seen when viewing the log file with `mysqlbinlog --base64-output=DECODE-ROWS` or in the output from `SHOW BINLOG EVENTS`.

As written in the output of statements such as `SHOW MASTER STATUS` or `SHOW SLAVE STATUS`, a sequence of GTIDs originating from the same server may be collapsed into a single expression, as shown here.

```
3E11FA47-71CA-11E1-9E33-C80AA9429562:1-5
```

The example just shown represents the first through fifth transactions originating on the MySQL Server whose `server_uuid` is `3E11FA47-71CA-11E1-9E33-C80AA9429562`.

This format is also used to supply the argument required by the `START SLAVE` options `SQL_BEFORE_GTIDS` and `SQL_AFTER_GTIDS`.

GTID Sets

A GTID set is a set of global transaction identifiers which is represented as shown here:

```
gtid_set:
    uuid_set [, uuid_set] ...
    | ''

uuid_set:
    uuid:interval[:interval]...

uuid:
    hhhhhhhh-hhhh-hhhh-hhhh-hhhhhhhhhhhh

h:
```

```
    [0-9|A-F]

interval:
    n[-n]

    (n >= 1)
```

GTID sets are used in the MySQL Server in several ways. For example, the values stored by the `gtid_executed` and `gtid_purged` system variables are represented as GTID sets. In addition, the functions `GTID_SUBSET()` and `GTID_SUBTRACT()` require GTID sets as input. When GTID sets are returned from server variables, UUIDs are in alphabetical order and numeric intervals are merged and in ascending order.

GTIDs are always preserved between master and slave. This means that you can always determine the source for any transaction applied on any slave by examining its binary log. In addition, once a transaction with a given GTID is committed on a given server, any subsequent transaction having the same GTID is ignored by that server. Thus, a transaction committed on the master can be applied no more than once on the slave, which helps to guarantee consistency.

When GTIDs are in use, the slave has no need for any nonlocal data, such as the name of a file on the master and a position within that file. All necessary information for synchronizing with the master is obtained directly from the replication data stream. GTIDs replace the file-offset pairs previously required to determine points for starting, stopping, or resuming the flow of data between master and slave. therefore, do not include `MASTER_LOG_FILE` or `MASTER_LOG_POS` options in the `CHANGE MASTER TO` statement used to direct a slave to replicate from a given master; instead it is necessary only to enable the `MASTER_AUTO_POSITION` option. For the exact steps needed to configure and start masters and slaves using GTID-based replication, see Section 18.1.3.2, “Setting Up Replication Using GTIDs”.

The generation and life cycle of a GTID consist of the following steps:

1. A transaction is executed and committed on the master.

   This transaction is assigned a GTID using the master's UUID and the smallest nonzero transaction sequence number not yet used on this server; the GTID is written to the master's binary log (immediately preceding the transaction itself in the log).

2. After the binary log data is transmitted to the slave and stored in the slave's relay log (using established mechanisms for this process—see Section 18.2, “Replication Implementation”, for details), the slave reads the GTID and sets the value of its `gtid_next` system variable as this GTID. This tells the slave that the next transaction must be logged using this GTID.

   It is important to note that the slave sets `gtid_next` in a session context.

3. The slave verifies that this GTID has not already been used to log a transaction in its own binary log. If this GTID has not been used, the slave then writes the GTID, applies the transaction, and writes the transaction to its binary log. By reading and checking the transaction's GTID first, before processing the transaction itself, the slave guarantees not only that no previous transaction having this GTID has been applied on the slave, but also that no other session has already read this GTID but has not yet committed the associated transaction. In other words, multiple clients are not permitted to apply the same transaction concurrently.

4. Because `gtid_next` is not empty, the slave does not attempt to generate a GTID for this transaction but instead writes the GTID stored in this variable—that is, the GTID obtained from the master—immediately preceding the transaction in its binary log.

mysql.gtid_executed Table

GTIDs are stored in a table named `gtid_executed`, in the `mysql` database. A row in this table contains, for each GTID or set of GTIDs that it represents, the UUID of the originating server, and the starting and ending transaction IDs of the set; for a row referencing only a single GTID, these last two values are the same.

The `mysql.gtid_executed` table is created (if it does not already exist) when the MySQL Server is installed or upgraded, using a `CREATE TABLE` statement similar to that shown here:

```
CREATE TABLE gtid_executed (
    source_uuid CHAR(36) NOT NULL,
    interval_start BIGINT(20) NOT NULL,
    interval_end BIGINT(20) NOT NULL,
    PRIMARY KEY (source_uuid, interval_start)
)
```

Warning

As with other MySQL system tables, do not attempt to create or modify this table yourself.

GTIDs are stored in the `mysql.gtid_executed` table only when `gtid_mode` is `ON` or `ON_PERMISSIVE`. GTIDs are stored in this table without regard to whether binary logging is enabled. However, the manner in which they are stored differs depending on whether `log_bin` is `ON` or `OFF`:

- If binary logging is disabled (`log_bin` is `OFF`), the server stores the GTID belonging to each transaction together with the transaction in the table.

  In addition, when binary logging is disabled, this table is compressed periodically at a user-configurable rate; see mysql.gtid_executed Table Compression, for more information.

- If binary logging is enabled (`log_bin` is `ON`), then in addition to storing the GTIDs in `mysql.gtid_executed`, whenever the binary log is rotated or the server is shut down, the server writes GTIDs for all transactions that were written into the previous binary log into the new binary log.

  In the event of the server stopping unexpectedly, the set of GTIDs from the previous binary log is not saved in the `mysql.gtid_executed` table. In this case, these GTIDs are added to the table and to the set of GTIDs in the `gtid_executed` system variable during recovery.

The `mysql.gtid_executed` table is reset by `RESET MASTER`.

mysql.gtid_executed Table Compression

Over the course of time, the `mysql.gtid_executed` table can become filled with many rows referring to individual GTIDs that originate on the same server, and whose transaction IDs make up a sequence, similar to what is shown here:

```
mysql> SELECT * FROM mysql.gtid_executed;
+--------------------------------------+----------------+--------------+
| source_uuid                          | interval_start | interval_end |
|--------------------------------------+----------------+--------------|
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 37             | 37           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 38             | 38           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 39             | 39           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 40             | 40           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 41             | 41           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 42             | 42           |
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 43             | 43           |
...
```

Considerable space can be saved if this table is compressed periodically by replacing each such set of rows with a single row that spans the entire interval of transaction identifiers, like this:

```
+--------------------------------------+----------------+--------------+
| source_uuid                          | interval_start | interval_end |
|--------------------------------------+----------------+--------------|
| 3E11FA47-71CA-11E1-9E33-C80AA9429562 | 37             | 43           |
```

```
...
```

When GTIDs are enabled, the server performs this type of compression on the `mysql.gtid_executed` table periodically. You can control the number of transactions that are allowed to elapse before the table is compressed, and thus the compression rate, by setting the `executed_gtids_compression_period` system variable. This variable's default value is 1000; this means that, by default, compression of the table is performed after each 1000 transactions. Setting `executed_gtid_compression_period` to 0 prevents the compression from being performed at all; however, you should be prepared for a potentially large increase in the amount of disk space that may be required by the `gtid_executed` table if you do this.

Note

When binary logging is enabled, the value of `executed_gtids_compression_period` is *not* used and the `mysql.gtid_executed` table is compressed on each binary log rotation.

Compression of the `mysql.gtid_executed` table is performed by a dedicated foreground thread that is created whenever GTIDs are enabled on the server. This thread is not listed in the output of `SHOW PROCESSLIST`, but it can be viewed as a row in the `threads` table, as shown here:

```
mysql> SELECT * FROM performance_schema.threads WHERE NAME LIKE '%gtid%'\G
*************************** 1. row ***************************
          THREAD_ID: 26
               NAME: thread/sql/compress_gtid_table
               TYPE: FOREGROUND
     PROCESSLIST_ID: 1
   PROCESSLIST_USER: NULL
   PROCESSLIST_HOST: NULL
     PROCESSLIST_DB: NULL
PROCESSLIST_COMMAND: Daemon
   PROCESSLIST_TIME: 1509
  PROCESSLIST_STATE: Suspending
   PROCESSLIST_INFO: NULL
   PARENT_THREAD_ID: 1
               ROLE: NULL
       INSTRUMENTED: YES
            HISTORY: YES
    CONNECTION_TYPE: NULL
       THREAD_OS_ID: 18677
```

This thread has the name `thread/sql/compress_gtid_table`, and normally sleeps until `executed_gtids_compression_period` transactions have been executed, then wakes up to perform compression of the `mysql.gtid_executed` table as described previously. It then sleeps until another `executed_gtids_compression_period` transactions have taken place, then wakes up to perform the compression again, repeating this loop indefinitely. Setting this value to 0 when binary logging is disabled means that the thread always sleeps and never wakes up.

18.1.3.2 Setting Up Replication Using GTIDs

This section describes a process for configuring and starting GTID-based replication in MySQL 8.0. This is a “cold start” procedure that assumes either that you are starting the replication master for the first time, or that it is possible to stop it; for information about provisioning replication slaves using GTIDs from a running master, see Section 18.1.3.3, “Using GTIDs for Failover and Scaleout”. For information about changing GTID mode on servers online, see Section 18.1.5, “Changing Replication Modes on Online Servers”.

The key steps in this startup process for the simplest possible GTID replication topology—consisting of one master and one slave—are as follows:

1. If replication is already running, synchronize both servers by making them read-only.
2. Stop both servers.

3. Restart both servers with GTIDs enabled and the correct options configured.

   The `mysqld` options necessary to start the servers as described are discussed in the example that follows later in this section.

   Note

   `server_uuid` must exist for GTIDs to function correctly.

4. Instruct the slave to use the master as the replication data source and to use auto-positioning, and then start the slave.

   The SQL statements needed to accomplish this step are described in the example that follows later in this section.

5. Enable read mode again on both servers, so that they can accept updates.

In the following example, two servers are already running as master and slave, using MySQL's binary log position-based replication protocol. If you are starting with new servers, see Section 18.1.2.3, "Creating a User for Replication" for information about adding a specific user for replication connections and Section 18.1.2.1, "Setting the Replication Master Configuration" for information about setting the server-id. The following examples show how to use startup options when running `mysqld`. Alternatively you can store startup options in an option file, see Section 4.2.6, "Using Option Files" for more information.

Most of the steps that follow require the use of the MySQL `root` account or another MySQL user account that has the `SUPER` privilege. `mysqladmin shutdown` requires either the `SUPER` privilege or the `SHUTDOWN` privilege.

Step 1: Synchronize the servers. Make the servers read-only. To do this, enable the `read_only` system variable by executing the following statement on both servers:

```
mysql> SET @@global.read_only = ON;
```

Then, allow the slave to catch up with the master. *It is extremely important that you make sure the slave has processed all updates before continuing*.

Step 2: Stop both servers. Stop each server using `mysqladmin` as shown here, where *`username`* is the user name for a MySQL user having sufficient privileges to shut down the server:

```
shell> mysqladmin -uusername -p shutdown
```

Then supply this user's password at the prompt.

Step 3: Restart both servers with GTIDs enabled. To enable GTID-based replication, each server must be started with GTID mode enabled, by setting the `--gtid-mode` option to `ON`, and with the `--enforce-gtid-consistency` option enabled to ensure that only statements which are safe for GTID-based replication are logged. In addition, you should start the slave with the `--skip-slave-start` option before configuring the slave settings. For more information on GTID related options, see Section 18.1.6.5, "Global Transaction ID Options and Variables".

It is not mandatory to have binary logging enabled in order to use GTIDs when using the mysql.gtid_executed Table. This means that you can have slave servers using GTIDs but without binary logging. Masters *must* always have binary logging enabled in order to be able to replicate. For example, to start a slave with GTIDs enabled but without binary logging, use at least these options:

```
shell> mysqld --gtid_mode=ON --enforce-gtid-consistency &
```

Depending on your configuration, supply additional options to `mysqld`.

Step 4: Direct the slave to use the master. Tell the slave to use the master as the replication data source, and to use GTID-based auto-positioning rather than file-based positioning. Execute a `CHANGE MASTER TO` statement on the slave, using the `MASTER_AUTO_POSITION` option to tell the slave that transactions will be identified by GTIDs.

You may also need to supply appropriate values for the master's host name and port number as well as the user name and password for a replication user account which can be used by the slave to connect to the master; if these have already been set prior to Step 1 and no further changes need to be made, the corresponding options can safely be omitted from the statement shown here.

```
mysql> CHANGE MASTER TO
     >     MASTER_HOST = host,
     >     MASTER_PORT = port,
     >     MASTER_USER = user,
     >     MASTER_PASSWORD = password,
     >     MASTER_AUTO_POSITION = 1;
```

Neither the `MASTER_LOG_FILE` option nor the `MASTER_LOG_POS` option may be used with `MASTER_AUTO_POSITION` set equal to 1. Attempting to do so causes the `CHANGE MASTER TO` statement to fail with an error.

Assuming that the `CHANGE MASTER TO` statement has succeeded, you can then start the slave, like this:

```
mysql> START SLAVE;
```

Step 5: Disable read-only mode. Allow the master to begin accepting updates once again by running the following statement:

```
mysql> SET @@global.read_only = OFF;
```

GTID-based replication should now be running, and you can begin (or resume) activity on the master as before. Section 18.1.3.3, “Using GTIDs for Failover and Scaleout”, discusses creation of new slaves when using GTIDs.

18.1.3.3 Using GTIDs for Failover and Scaleout

There are a number of techniques when using MySQL Replication with Global Transaction Identifiers (GTIDs) for provisioning a new slave which can then be used for scaleout, being promoted to master as necessary for failover. This section describes the following techniques:

- Simple replication
- Copying data and transactions to the slave
- Injecting empty transactions
- Excluding transactions with gtid_purged
- Restoring GTID mode slaves

Global transaction identifiers were added to MySQL Replication for the purpose of simplifying in general management of the replication data flow and of failover activities in particular. Each identifier uniquely identifies a set of binary log events that together make up a transaction. GTIDs play a key role in applying changes to the database: the server automatically skips any transaction having an identifier which the server recognizes as one that it has processed before. This behavior is critical for automatic replication positioning and correct failover.

The mapping between identifiers and sets of events comprising a given transaction is captured in the binary log. This poses some challenges when provisioning a new server with data from another existing server. To reproduce the identifier set on the new server, it is necessary to copy the identifiers from the old server to the new one, and to preserve the relationship between the identifiers and the actual

events. This is neccessary for restoring a slave that is immediately available as a candidate to become a new master on failover or switchover.

Simple replication. The easiest way to reproduce all identifiers and transactions on a new server is to make the new server into the slave of a master that has the entire execution history, and enable global transaction identifiers on both servers. See Section 18.1.3.2, “Setting Up Replication Using GTIDs”, for more information.

Once replication is started, the new server copies the entire binary log from the master and thus obtains all information about all GTIDs.

This method is simple and effective, but requires the slave to read the binary log from the master; it can sometimes take a comparatively long time for the new slave to catch up with the master, so this method is not suitable for fast failover or restoring from backup. This section explains how to avoid fetching all of the execution history from the master by copying binary log files to the new server.

Copying data and transactions to the slave. Playing back the entire transaction history can be time-consuming, and represents a major bottleneck when setting up a new replication slave. To eliminate this requirement, a snapshot of the data set, the binary logs and the global transaction information the master contains is imported to the slave. The binary log is played back, after which replication can be started, allowing the slave to become current with any remaining transactions.

There are several variants of this method, the difference being in the manner in which data dumps and transactions from binary logs are transfered to the slave, as outlined here:

Data Set	Transaction History

• Use the `mysql` client to import a dump file created with `mysqldump`. Use the `--master-data` option to include binary logging information and `--set-gtid-purged` to `AUTO` (the default) or `ON`, to include information about executed transactions. You should have `--gtid-mode=ON` while importing the dump on the slave.
• Stop the slave, copy the contents of the master's data directory to the slave's data directory, then restart the slave.
	If `gtid_mode` is not `ON`, restart the server with GTID mode enabled.
	• Import the binary log using `mysqlbinlog`, with the `--read-from-remote-server` and `--read-from-remote-master` options.
	• Copy the master's binary log files to the slave. You can make copies from the slave using `mysqlbinlog --read-from-remote-server --raw`. These can be read in to the slave in either of the following ways:
	  • Update the slave's `binlog.index` file to point to the copied log files. Then execute a `CHANGE MASTER TO` statement in the `mysql` client to point to the first log file, and `START SLAVE` to read them.
	  • Use `mysqlbinlog > file` (without the `--raw` option) to export the binary log files to SQL files that can be processed by the `mysql` client.

See also Section 4.6.8.3, “Using mysqlbinlog to Back Up Binary Log Files”.

This method has the advantage that a new server is available almost immediately; only those transactions that were committed while the snapshot or dump file was being replayed still need to be obtained from the existing master. This means that the slave's availability is not instantanteous—but only a relatively short amount of time should be required for the slave to catch up with these few remaining transactions.

Copying over binary logs to the target server in advance is usually faster than reading the entire transaction execution history from the master in real time. However, it may not always be feasible to

move these files to the target when required, due to size or other considerations. The two remaining methods for provisioning a new slave discussed in this section use other means to transfer information about transactions to the new slave.

Injecting empty transactions. The master's global `gtid_executed` variable contains the set of all transactions executed on the master. Rather than copy the binary logs when taking a snapshot to provision a new server, you can instead note the content of `gtid_executed` on the server from which the snapshot was taken. Before adding the new server to the replication chain, simply commit an empty transaction on the new server for each transaction identifier contained in the master's `gtid_executed`, like this:

```
SET GTID_NEXT='aaa-bbb-ccc-ddd:N';

BEGIN;
COMMIT;

SET GTID_NEXT='AUTOMATIC';
```

Once all transaction identifiers have been reinstated in this way using empty transactions, you must flush and purge the slave's binary logs, as shown here, where *`N`* is the nonzero suffix of the current binary log file name:

```
FLUSH LOGS;
PURGE BINARY LOGS TO 'master-bin.00000N';
```

You should do this to prevent this server from flooding the replication stream with false transactions in the event that it is later promoted to master. (The `FLUSH LOGS` statement forces the creation of a new binary log file; `PURGE BINARY LOGS` purges the empty transactions, but retains their identifiers.)

This method creates a server that is essentially a snapshot, but in time is able to become a master as its binary log history converges with that of the replication stream (that is, as it catches up with the master or masters). This outcome is similar in effect to that obtained using the remaining provisioning method, which we discuss in the next few paragraphs.

Excluding transactions with gtid_purged. The master's global `gtid_purged` variable contains the set of all transactions that have been purged from the master's binary log. As with the method discussed previously (see Injecting empty transactions), you can record the value of `gtid_executed` on the server from which the snapshot was taken (in place of copying the binary logs to the new server). Unlike the previous method, there is no need to commit empty transactions (or to issue `PURGE BINARY LOGS`); instead, you can set `gtid_purged` on the slave directly, based on the value of `gtid_executed` on the server from which the backup or snapshot was taken.

As with the method using empty transactions, this method creates a server that is functionally a snapshot, but in time is able to become a master as its binary log history converges with that of the replication master or group.

Restoring GTID mode slaves. When restoring a slave in a GTID based replication setup that has encountered an error, injecting an empty transaction may not solve the problem because an event does not have a GTID.

Use `mysqlbinlog` to find the next transaction, which is probably the first transaction in the next log file after the event. Copy everything up to the `COMMIT` for that transaction, being sure to include the `SET @@SESSION.GTID_NEXT`. Even if you are not using row-based replication, you can still run binary log row events in the command line client.

Stop the slave and run the transaction you copied. The `mysqlbinlog` output sets the delimiter to `/*! */;`, so set it back:

```
mysql> DELIMITER ;
```

Restart replication from the correct position automatically:

```
mysql> SET GTID_NEXT=automatic;
mysql> RESET SLAVE;
mysql> START SLAVE;
```

18.1.3.4 Restrictions on Replication with GTIDs

Because GTID-based replication is dependent on transactions, some features otherwise available in MySQL are not supported when using it. This section provides information about restrictions on and limitations of replication with GTIDs.

Updates involving nontransactional storage engines. When using GTIDs, updates to tables using nontransactional storage engines such as `MyISAM` cannot be made in the same statement or transaction as updates to tables using transactional storage engines such as `InnoDB`.

This restriction is due to the fact that updates to tables that use a nontransactional storage engine mixed with updates to tables that use a transactional storage engine within the same transaction can result in multiple GTIDs being assigned to the same transaction.

Such problems can also occur when the master and the slave use different storage engines for their respective versions of the same table, where one storage engine is transactional and the other is not.

In any of the cases just mentioned, the one-to-one correspondence between transactions and GTIDs is broken, with the result that GTID-based replication cannot function correctly.

CREATE TABLE ... SELECT statements. `CREATE TABLE ... SELECT` is not safe for statement-based replication. When using row-based replication, this statement is actually logged as two separate events—one for the creation of the table, and another for the insertion of rows from the source table into the new table just created. When this statement is executed within a transaction, it is possible in some cases for these two events to receive the same transaction identifier, which means that the transaction containing the inserts is skipped by the slave. Therefore, `CREATE TABLE ... SELECT` is not supported when using GTID-based replication.

Temporary tables. `CREATE TEMPORARY TABLE` and `DROP TEMPORARY TABLE` statements are not supported inside transactions when using GTIDs (that is, when the server was started with the `--enforce-gtid-consistency` option). It is possible to use these statements with GTIDs enabled, but only outside of any transaction, and only with `autocommit=1`.

Preventing execution of unsupported statements. To prevent execution of statements that would cause GTID-based replication to fail, all servers must be started with the `--enforce-gtid-consistency` option when enabling GTIDs. This causes statements of any of the types discussed previously in this section to fail with an error.

For information about other required startup options when enabling GTIDs, see Section 18.1.3.2, “Setting Up Replication Using GTIDs”.

`sql_slave_skip_counter` is not supported when using GTIDs. If you need to skip transactions, use the value of the master's `gtid_executed` variable instead; see Injecting empty transactions, for more information.

GTID mode and mysqldump. It is possible to import a dump made using `mysqldump` into a MySQL Server running with GTID mode enabled, provided that there are no GTIDs in the target server's binary log.

GTID mode and mysql_upgrade. It is not recommended to use `mysql_upgrade` with the `--write-binlog` option on a MySQL Server running with `--gtid-mode=ON` because `mysql_upgrade` can make changes to system tables that use the `MyISAM` storage engine, which is non-transactional.

18.1.4 MySQL Multi-Source Replication

This section describes MySQL Multi-Source Replication, which enables you to replicate from multiple immediate masters in parallel. This section describes multi-source replication, and how to configure, monitor and troubleshoot it.

18.1.4.1 MySQL Multi-Source Replication Overview

MySQL Multi-Source Replication enables a replication slave to receive transactions from multiple sources simultaneously. Multi-source replication can be used to back up multiple servers to a single server, to merge table shards, and consolidate data from multiple servers to a single server. Multi-source replication does not implement any conflict detection or resolution when applying the transactions, and those tasks are left to the application if required. In a multi-source replication topology, a slave creates a replication channel for each master that it should receive transactions from. See Section 18.2.3, “Replication Channels”. The following sections describe how to set up multi-source replication.

18.1.4.2 Multi-Source Replication Tutorials

This section provides tutorials on how to configure masters and slaves for multi-source replication, and how to start, stop and reset multi-source slaves.

Configuring Multi-Source Replication

This section explains how to configure a multi-source replication topology, and provides details about configuring masters and slaves. Such a topology requires at least two masters and one slave configured.

Masters in a multi-source replication topology can be configured to use either global transaction identifier (GTID) based replication, or binary log position-based replication. See Section 18.1.3.2, “Setting Up Replication Using GTIDs” for how to configure a master using GTID based replication. See Section 18.1.2.1, “Setting the Replication Master Configuration” for how to configure a master using file position based replication.

Slaves in a multi-source replication topology require `TABLE` based repositories. Multi-source replication is not compatible with `FILE` based repositories. The type of repository being used by `mysqld` can be configured either at startup, or dynamically.

To configure the type of repository used by a replication slave at startup, start `mysqld` with the following options:

```
--master-info-repository=TABLE --relay-log-info-repository=TABLE
```

To modify an existing replication slave that is using a `FILE` repository to use `TABLE` repositories, convert the existing replication repositories dynamically by running the following commands:

```
STOP SLAVE;
SET GLOBAL master_info_repository = 'TABLE';
SET GLOBAL relay_log_info_repository = 'TABLE';
```

Adding a GTID Based Master to a Multi-Source Replication Slave

This section assumes you have enabled GTID based transactions on the master using `gtid_mode=ON`, enabled a replication user, and ensured that the slave is using `TABLE` based replication repositories. Use the `CHANGE MASTER TO` statement to add a new master to a channel by using a `FOR CHANNEL` *`channel`* clause. For more information on replication channels, see Section 18.2.3, “Replication Channels”

For example, to add a new master with the host name *`master1`* using port *`3451`* to a channel called *`master-1`*:

```
CHANGE MASTER TO MASTER_HOST='master1', MASTER_USER='rpl', MASTER_PORT=3451, MASTER_PASSWORD='', \
MASTER_AUTO_POSITION = 1 FOR CHANNEL 'master-1';
```

Multi-source replication is compatible with auto-positioning. See Section 13.4.2.1, “CHANGE MASTER TO Syntax” for more information.

Repeat this process for each extra master that you want to add to a channel, changing the host name, port and channel as appropriate.

Adding a Binary Log Based Master to a Multi-Source Replication Slave

This section assumes you have enabled binary logging on the master using `--log-bin`, enabled a replication user, noted the current binary log position, and ensured that the slave is using `TABLE` based replication repositories. You need to know the current `MASTER_LOG_FILE` and `MASTER_LOG_POSITION`. Use the `CHANGE MASTER TO` statement to add a new master to a channel by specifying a `FOR CHANNEL channel` clause. For example, to add a new master with the host name `master1` using port `3451` to a channel called `master-1`:

```
CHANGE MASTER TO MASTER_HOST='master1', MASTER_USER='rpl', MASTER_PORT=3451, MASTER_PASSWORD='' \
MASTER_LOG_FILE='master1-bin.000006', MASTER_LOG_POS=628 FOR CHANNEL 'master-1';
```

Repeat this process for each extra master that you want to add to a channel, changing the host name, port and channel as appropriate.

Starting Multi-Source Replication Slaves

Once you have added all of the channels you want to use as replication masters, use a `START SLAVE thread_types` statement to start replication. When you have enabled multiple channels on a slave, you can choose to either start all channels, or select a specific channel to start.

- To start all currently configured replication channels:

```
START SLAVE thread_types;
```

- To start only a named channel, use a `FOR CHANNEL channel` clause:

```
START SLAVE thread_types FOR CHANNEL channel;
```

Use the `thread_types` option to choose specific threads you want the above statements to start on the slave. See Section 13.4.2.6, “START SLAVE Syntax” for more information.

Stopping Multi-Source Replication Slaves

The `STOP SLAVE` statement can be used to stop a multi-source replication slave. By default, if you use the `STOP SLAVE` statement on a multi-source replication slave all channels are stopped. Optionally, use the `FOR CHANNEL channel` clause to stop only a specific channel.

- To stop all currently configured replication channels:

```
STOP SLAVE thread_types;
```

- To stop only a named channel, use a `FOR CHANNEL channel` clause:

```
STOP SLAVE thread_types FOR CHANNEL channel;
```

Use the `thread_types` option to choose specific threads you want the above statements to stop on the slave. See Section 13.4.2.7, “STOP SLAVE Syntax” for more information.

Resetting Multi-Source Replication Slaves

The `RESET SLAVE` statement can be used to reset a multi-source replication slave. By default, if you use the `RESET SLAVE` statement on a multi-source replication slave all channels are reset. Optionally, use the `FOR CHANNEL channel` clause to reset only a specific channel.

- To reset all currently configured replication channels:

```
RESET SLAVE;
```

- To reset only a named channel, use a `FOR CHANNEL channel` clause:

```
RESET SLAVE FOR CHANNEL channel;
```

See Section 13.4.2.4, “RESET SLAVE Syntax” for more information.

18.1.4.3 Multi-Source Replication Monitoring

To monitor the status of replication channels the following options exist:

- Using the replication Performance Schema tables. The first column of these tables is `Channel_Name`. This enables you to write complex queries based on `Channel_Name` as a key. See Section 25.11.11, “Performance Schema Replication Tables”.
- Using `SHOW SLAVE STATUS FOR CHANNEL channel_name`. By default, if the `FOR CHANNEL channel_name` clause is not used, this command shows the slave status for all channels with one row per channel. The identifier `channel_name` is added as a column in the result set. If a `FOR CHANNEL channel_name` clause is provided, the results show the status of only the named replication channel.

Note

The `SHOW VARIABLES` statement does not work with multiple replication channels. The information that was available through these variables has been migrated to the replication performance tables. Using a `SHOW VARIABLES` statement in a topology with multiple channels shows the status of only the default channel.

Monitoring Channels Using Performance Schema Tables

This section explains how to use the replication Performance Schema tables to monitor channels. You can choose to monitor all channels, or a subset of the existing channels.

To monitor the connection status of all channels:

```
mysql> SELECT * FROM replication_connection_status\G;
*************************** 1. row ***************************
CHANNEL_NAME: master1
GROUP_NAME:
SOURCE_UUID: 046e41f8-a223-11e4-a975-0811960cc264
THREAD_ID: 24
SERVICE_STATE: ON
COUNT_RECEIVED_HEARTBEATS: 0
LAST_HEARTBEAT_TIMESTAMP: 0000-00-00 00:00:00
RECEIVED_TRANSACTION_SET: 046e41f8-a223-11e4-a975-0811960cc264:4-37
LAST_ERROR_NUMBER: 0
LAST_ERROR_MESSAGE:
LAST_ERROR_TIMESTAMP: 0000-00-00 00:00:00
*************************** 2. row ***************************
CHANNEL_NAME: master2
GROUP_NAME:
SOURCE_UUID: 7475e474-a223-11e4-a978-0811960cc264
THREAD_ID: 26
SERVICE_STATE: ON
COUNT_RECEIVED_HEARTBEATS: 0
LAST_HEARTBEAT_TIMESTAMP: 0000-00-00 00:00:00
RECEIVED_TRANSACTION_SET: 7475e474-a223-11e4-a978-0811960cc264:4-6
LAST_ERROR_NUMBER: 0
LAST_ERROR_MESSAGE:
LAST_ERROR_TIMESTAMP: 0000-00-00 00:00:00
```

```
2 rows in set (0.00 sec)
```

In the above output there are two channels enabled, and as shown by the `CHANNEL_NAME` field they are called `master1` and `master2`.

The addition of the `CHANNEL_NAME` field enables you to query the Performance Schema tables for a specific channel. To monitor the connection status of a named channel, use a `WHERE channel_name=channel` clause:

```
mysql> SELECT * FROM replication_connection_status WHERE channel_name='master1'\G
*************************** 1. row ***************************
CHANNEL_NAME: master1
GROUP_NAME:
SOURCE_UUID: 046e41f8-a223-11e4-a975-0811960cc264
THREAD_ID: 24
SERVICE_STATE: ON
COUNT_RECEIVED_HEARTBEATS: 0
LAST_HEARTBEAT_TIMESTAMP: 0000-00-00 00:00:00
RECEIVED_TRANSACTION_SET: 046e41f8-a223-11e4-a975-0811960cc264:4-37
LAST_ERROR_NUMBER: 0
LAST_ERROR_MESSAGE:
LAST_ERROR_TIMESTAMP: 0000-00-00 00:00:00
1 row in set (0.00 sec)
```

Similarly, the `WHERE channel_name=channel` clause can be used to monitor the other replication Performance Schema tables for a specific channel. For more information, see Section 25.11.11, “Performance Schema Replication Tables”.

18.1.4.4 Multi-Source Replication Error Messages

Error codes and messages provide information about errors encountered in a multi-source replication topology. These error codes and messages are only emitted when multi-source replication is enabled, and provide information related to the channel which generated the error. For example:

`Slave is already running` and `Slave is already stopped` have been replaced with `Replication thread(s) for channel channel_name are already running` and `Replication threads(s) for channel channel_name are already stopped` respectively.

The server log messages have also been changed to indicate which channel the log messages relate to. This makes debugging and tracing easier.

18.1.5 Changing Replication Modes on Online Servers

This section describes how to change the mode of replication being used without having to take the server offline.

18.1.5.1 Replication Mode Concepts

To be able to safely configure the replication mode of an online server it is important to understand some key concepts of replication. This section explains these concepts and is essential reading before attempting to modify the replication mode of an online server.

The modes of replication available in MySQL rely on different techniques for identifying transactions which are logged. The types of transactions used by replication are as follows:

- GTID transactions are identified by a global transaction identifier (GTID) in the form `UUID:NUMBER`. Every GTID transaction in a log is always preceded by a `Gtid_log_event`. GTID transactions can be addressed using either the GTID or using the file name and position.
- Anonymous transactions do not have a GTID assigned, and MySQL ensures that every anonymous transaction in a log is preceded by an `Anonymous_gtid_log_event`. In previous versions,

anonymous transactions were not preceded by any particular event. Anonymous transactions can only be addressed using file name and position.

When using GTIDs you can take advantage of auto-positioning and automatic fail-over, as well as use `WAIT_FOR_EXECUTED_GTID_SET()`, `session_track_gtids`, and monitor replicated transactions using Performance Schema tables. With GTIDs enabled you cannot use `sql_slave_skip_counter`, instead use empty transactions.

Transactions in a relay log that was received from a master running a previous version of MySQL may not be preceded by any particular event at all, but after being replayed and logged in the slave's binary log, they are preceded with an `Anonymous_gtid_log_event`.

The ability to configure the replication mode online means that the `gtid_mode` and `enforce_gtid_consistency` variables are now both dynamic and can be set by `SUPER` from a top-level statement. In MySQL 5.6 and below, both of these variables could only be configured using the appropriate option at server start, meaning that changes to the replication mode required a server restart. In all versions `gtid_mode` could be set to `ON` or `OFF`, which corresponded to whether GTIDs were used to identify transactions or not. When `gtid_mode=ON` it is not possible to replicate anonymous transactions, and when `gtid_mode=OFF` only anonymous transactions can be replicated. When `gtid_mode=OFF_PERMISSIVE` then *new* transactions are anonymous while permitting replicated transactions to be either GTID or anonymous transactions. When `gtid_mode=ON_PERMISSIVE` then *new* transactions use GTIDs while permitting replicated transactions to be either GTID or anonymous transactions. This means it is possible to have a replication topology that has servers using both anonymous and GTID transactions. For example a master with `gtid_mode=ON` could be replicating to a slave with `gtid_mode=ON_PERMISSIVE`. The valid values for `gtid_mode` are as follows and in this order:

- `OFF`
- `OFF_PERMISSIVE`
- `ON_PERMISSIVE`
- `ON`

It is important to note that the state of `gtid_mode` can only be changed by one step at a time based on the above order. For example, if `gtid_mode` is currently set to `OFF_PERMISSIVE`, it is possible to change to `OFF` or `ON_PERMISSIVE` but not to `ON`. This is to ensure that the process of changing from anonymous transactions to GTID transactions online is correctly handled by the server. When you switch between `gtid_mode=ON` and `gtid_mode=OFF`, the GTID state (in other words the value of `gtid_executed`) is persistent. This ensures that the GTID set that has been applied by the server is always retained, regardless of changes between types of `gtid_mode`.

The fields related to GTIDs display the correct information regardless of the currently selected `gtid_mode`. This means that fields which display GTID sets, such as `gtid_executed`, `gtid_purged`, `RECEIVED_TRANSACTION_SET` in the `replication_connection_status` Performance Schema table, and the GTID related results of `SHOW SLAVE STATUS`, now return the empty string when there are no GTIDs present. Fields that display a single GTID, such as `CURRENT_TRANSACTION` in the `replication_applier_status_by_worker` Performance Schema table, now display `ANONYMOUS` when GTID transactions are not being used.

Replication from a master using `gtid_mode=ON` provides the ability to use auto-positioning, configured using the `CHANGE MASTER TO MASTER_AUTO_POSITION = 1;` statement. The replication topology being used impacts on whether it is possible to enable auto-positioning or not, as this feature relies on GTIDs and is not compatible with anonymous transactions. An error is generated if auto-positioning is enabled and an anonymous transaction is encountered. It is strongly recommended to ensure there are no anonymous transactions remaining in the topology before enabling auto-positioning, see Section 18.1.5.2, “Enabling GTID Transactions Online”. The valid combinations of `gtid_mode` and auto-positioning on master and slave are shown in the following table, where the master's `gtid_mode` is shown on the horizontal and the slave's `gtid_mode` is on the vertical:

Table 18.1 Valid Combinations of Master and Slave gtid_mode

Master/Slave gtid_mode	OFF	OFF_PERMISSIVE	ON_PERMISSIVE	ON
OFF	Y	Y	N	N
OFF_PERMISSIVE	Y	Y	Y	Y*
ON_PERMISSIVE	Y	Y	Y	Y*
ON	N	N	Y	Y*

In the above table, the entries are:

- Y: the gtid_mode of master and slave is compatible
- N: the gtid_mode of master and slave is not compatible
- *: auto-positioning can be used

The currently selected gtid_mode also impacts on the gtid_next variable. The following table shows the behavior of the server for the different values of gtid_mode and gtid_next.

Table 18.2 Valid Combinations of gtid_mode and gtid_next

gtid_next	AUTOMATIC binary log on	AUTOMATIC binary log off	ANONYMOUS	UUID:NUMBER
OFF	ANONYMOUS	ANONYMOUS	ANONYMOUS	Error
OFF_PERMISSIVE	ANONYMOUS	ANONYMOUS	ANONYMOUS	UUID:NUMBER
ON_PERMISSIVE	New GTID	ANONYMOUS	ANONYMOUS	UUID:NUMBER
ON	New GTID	ANONYMOUS	Error	UUID:NUMBER

In the above table, the entries are:

- ANONYMOUS: generate an anonymous transaction.
- Error: generate an error and fail to execute SET GTID_NEXT.
- UUID:NUMBER: generate a GTID with the specified UUID:NUMBER.
- New GTID: generate a GTID with an automatically generated number.

When the binary log is off and gtid_next is set to AUTOMATIC, then no GTID is generated. This is consistent with the behavior of previous versions.

18.1.5.2 Enabling GTID Transactions Online

This section describes how to enable GTID transactions, and optionally auto-positioning, on servers that are already online and using anonymous transactions. This procedure does not require taking the server offline and is suited to use in production. However, if you have the possibility to take the servers offline when enabling GTID transactions that process is easier.

Before you start, ensure that the servers meet the following pre-conditions:

- *All* servers in your topology must use MySQL 5.7.6 or later. You cannot enable GTID transactions online on any single server unless *all* servers which are in the topology are using this version.
- All servers have gtid_mode set to the default value OFF.

The following procedure can be paused at any time and later resumed where it was, or reversed by jumping to the corresponding step of Section 18.1.5.3, “Disabling GTID Transactions Online”, the

online procedure to disable GTIDs. This makes the procedure fault-tolerant because any unrelated issues that may appear in the middle of the procedure can be handled as usual, and then the procedure continued where it was left off.

Note

It is crucial that you complete every step before continuing to the next step.

To enable GTID transactions:

1. On each server, execute:

   ```
   SET @@GLOBAL.ENFORCE_GTID_CONSISTENCY = WARN;
   ```

   Let the server run for a while with your normal workload and monitor the logs. If this step causes any warnings in the log, adjust your application so that it only uses GTID-compatible features and does not generate any warnings.

   Important

   This is the first important step. You must ensure that no warnings are being generated in the error logs before going to the next step.

2. On each server, execute:

   ```
   SET @@GLOBAL.ENFORCE_GTID_CONSISTENCY = ON;
   ```

3. On each server, execute:

   ```
   SET @@GLOBAL.GTID_MODE = OFF_PERMISSIVE;
   ```

   It does not matter which server executes this statement first, but it is important that all servers complete this step before any server begins the next step.

4. On each server, execute:

   ```
   SET @@GLOBAL.GTID_MODE = ON_PERMISSIVE;
   ```

   It does not matter which server executes this statement first.

5. On each server, wait until the status variable `ONGOING_ANONYMOUS_TRANSACTION_COUNT` is zero. This can be checked using:

   ```
   SHOW STATUS LIKE 'ONGOING_ANONYMOUS_TRANSACTION_COUNT';
   ```

   Note

   On a replication slave, it is theoretically possible that this shows zero and then non-zero again. This is not a problem, it suffices that it shows zero once.

6. Wait for all transactions generated up to step 5 to replicate to all servers. You can do this without stopping updates: the only important thing is that all anonymous transactions get replicated.

   See Section 18.1.5.4, “Verifying Replication of Anonymous Transactions” for one method of checking that all anonymous transactions have replicated to all servers.

7. If you use binary logs for anything other than replication, for example point in time backup and restore, wait until you do not need the old binary logs having transactions without GTIDs.

   For instance, after step 6 has completed, you can execute `FLUSH LOGS` on the server where you are taking backups. Then either explicitly take a backup or wait for the next iteration of any periodic backup routine you may have set up.

Ideally, wait for the server to purge all binary logs that existed when step 6 was completed. Also wait for any backup taken before step 6 to expire.

Important

This is the second important point. It is vital to understand that binary logs containing anonymous transactions, without GTIDs cannot be used after the next step. After this step, you must be sure that transactions without GTIDs do not exist anywhere in the topology.

8. On each server, execute:

```
SET @@GLOBAL.GTID_MODE = ON;
```

9. On each server, add `gtid-mode=ON` to `my.cnf`.

   You are now guaranteed that all transactions have a GTID (except transactions generated in step 5 or earlier, which have already been processed). To start using the GTID protocol so that you can later perform automatic fail-over, execute the following on each slave. Optionally, if you use multi-source replication, do this for each channel and include the `FOR CHANNEL` *`channel`* clause:

```
STOP SLAVE [FOR CHANNEL 'channel'];
CHANGE MASTER TO MASTER_AUTO_POSITION = 1 [FOR CHANNEL 'channel'];
START SLAVE [FOR CHANNEL 'channel'];
```

18.1.5.3 Disabling GTID Transactions Online

This section describes how to disable GTID transactions on servers that are already online. This procedure does not require taking the server offline and is suited to use in production. However, if you have the possibility to take the servers offline when disabling GTIDs mode that process is easier.

The process is similar to enabling GTID transactions while the server is online, but reversing the steps. The only thing that differs is the point at which you wait for logged transactions to replicate.

Before you start, ensure that the servers meet the following pre-conditions:

- *All* servers in your topology must use MySQL 5.7.6 or later. You cannot disable GTID transactions online on any single server unless *all* servers which are in the topology are using this version.
- All servers have `gtid_mode` set to `ON`.

1. Execute the following on each slave, and if you using multi-source replication, do it for each channel and include the `FOR CHANNEL` channel clause:

```
STOP SLAVE [FOR CHANNEL 'channel'];
CHANGE MASTER TO MASTER_AUTO_POSITION = 0, MASTER_LOG_FILE = file, \
MASTER_LOG_POS = position [FOR CHANNEL 'channel'];
START SLAVE [FOR CHANNEL 'channel'];
```

2. On each server, execute:

```
SET @@GLOBAL.GTID_MODE = ON_PERMISSIVE;
```

3. On each server, execute:

```
SET @@GLOBAL.GTID_MODE = OFF_PERMISSIVE;
```

4. On each server, wait until the variable @@GLOBAL.GTID_OWNED is equal to the empty string. This can be checked using:

```
SELECT @@GLOBAL.GTID_OWNED;
```

On a replication slave, it is theoretically possible that this is empty and then nonempty again. This is not a problem, it suffices that it is empty once.

5. Wait for all transactions that currently exist in any binary log to replicate to all slaves. See Section 18.1.5.4, "Verifying Replication of Anonymous Transactions" for one method of checking that all anonymous transactions have replicated to all servers.

6. If you use binary logs for anything else than replication, for example to do point in time backup or restore: wait until you do not need the old binary logs having GTID transactions.

   For instance, after step 5 has completed, you can execute `FLUSH LOGS` on the server where you are taking the backup. Then either explicitly take a backup or wait for the next iteration of any periodic backup routine you may have set up.

   Ideally, wait for the server to purge all binary logs that existed when step 5 was completed. Also wait for any backup taken before step 5 to expire.

   Important

   This is the one important point during this procedure. It is important to understand that logs containing GTID transactions cannot be used after the next step. Before proceeding you must be sure that GTID transactions do not exist anywhere in the topology.

7. On each server, execute:

   ```
   SET @@GLOBAL.GTID_MODE = OFF;
   ```

8. On each server, set `gtid-mode=OFF` in `my.cnf`.

   If you want to set `enforce_gtid_consistency=OFF`, you can do so now. After setting it, you should add `enforce_gtid_consistency=OFF` to your configuration file.

If you want to downgrade to an earlier version of MySQL, you can do so now, using the normal downgrade procedure.

18.1.5.4 Verifying Replication of Anonymous Transactions

This section explains how to monitor a replication topology and verify that all anonymous transactions have been replicated. This is helpful when changing the replication mode online as you can verify that it is safe to change to GTID transactions.

There are several possible ways to wait for transactions to replicate:

The simplest method, which works regardless of your topology but relies on timing is as follows: if you are sure that the slave never lags more than N seconds, just wait for a bit more than N seconds. Or wait for a day, or whatever time period you consider safe for your deployment.

A safer method in the sense that it does not depend on timing: if you only have a master with one or more slaves, do the following:

1. On the master, execute:

   ```
   SHOW MASTER STATUS;
   ```

   Note down the values in the `File` and `Position` column.

2. On every slave, use the file and position information from the master to execute:

   ```
   SELECT MASTER_POS_WAIT(file, position);
   ```

If you have a master and multiple levels of slaves, or in other words you have slaves of slaves, repeat step 2 on each level, starting from the master, then all the direct slaves, then all the slaves of slaves, and so on.

If you use a circular replication topology where multiple servers may have write clients, perform step 2 for each master-slave connection, until you have completed the full circle. Repeat the whole process so that you do the full circle *twice*.

For example, suppose you have three servers A, B, and C, replicating in a circle so that A -> B -> C -> A. The procedure is then:

- Do step 1 on A and step 2 on B.
- Do step 1 on B and step 2 on C.
- Do step 1 on C and step 2 on A.
- Do step 1 on A and step 2 on B.
- Do step 1 on B and step 2 on C.
- Do step 1 on C and step 2 on A.

18.1.6 Replication and Binary Logging Options and Variables

The following sections contain information about `mysqld` options and server variables that are used in replication and for controlling the binary log. Options and variables for use on replication masters and replication slaves are covered separately, as are options and variables relating to binary logging and global transaction identifiers (GTIDs). A set of quick-reference tables providing basic information about these options and variables is also included.

Of particular importance is the `--server-id` option.

Command-Line Format	`--server-id=#`	
System Variable	**Name**	`server_id`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

This option is common to both master and slave replication servers, and is used in replication to enable master and slave servers to identify themselves uniquely. For additional information, see Section 18.1.6.2, “Replication Master Options and Variables”, and Section 18.1.6.3, “Replication Slave Options and Variables”.

On the master and each slave, you *must* use the `--server-id` option to establish a unique replication ID in the range from 1 to $2^{32} - 1$. “Unique”, means that each ID must be different from every other ID in use by any other replication master or slave. For example, `server-id=3`.

The `--server-id` must be used if binary logging is enabled, and a value of 0 is not changed by the server. If you specify `--server-id` without an argument, the effect is the same as using 0. In either case, if the `server_id` is 0, binary logging takes place, but slaves cannot connect to the master, nor can any other servers connect to it as slaves. (Bug #11763963, Bug #56718)

For more information, see Section 18.1.2.2, “Setting the Replication Slave Configuration”.

`server_uuid`

In MySQL 8.0, the server generates a true UUID in addition to the `--server-id` supplied by the user. This is available as the global, read-only variable `server_uuid`.

Note

The presence of the `server_uuid` system variable in MySQL 8.0 does not change the requirement for setting a unique `--server-id` for each MySQL server as part of preparing and running MySQL replication, as described earlier in this section.

System Variable	Name	`server_uuid`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

When starting, the MySQL server automatically obtains a UUID as follows:

1. Attempt to read and use the UUID written in the file *`data_dir`*`/auto.cnf` (where *`data_dir`* is the server's data directory).
2. If *`data_dir`*`/auto.cnf` is not found, generate a new UUID and save it to this file, creating the file if necessary.

The `auto.cnf` file has a format similar to that used for `my.cnf` or `my.ini` files. In MySQL 8.0, `auto.cnf` has only a single `[auto]` section containing a single `server_uuid` setting and value; the file's contents appear similar to what is shown here:

```
[auto]
server_uuid=8a94f357-aab4-11df-86ab-c80aa9429562
```

Important

The `auto.cnf` file is automatically generated; do not attempt to write or modify this file.

When using MySQL replication, masters and slaves know each other's UUIDs. The value of a slave's UUID can be seen in the output of `SHOW SLAVE HOSTS`. Once `START SLAVE` has been executed, the value of the master's UUID is available on the slave in the output of `SHOW SLAVE STATUS`.

Note

Issuing a `STOP SLAVE` or `RESET SLAVE` statement does *not* reset the master's UUID as used on the slave.

A server's `server_uuid` is also used in GTIDs for transactions originating on that server. For more information, see Section 18.1.3, “Replication with Global Transaction Identifiers”.

When starting, the slave I/O thread generates an error and aborts if its master's UUID is equal to its own unless the `--replicate-same-server-id` option has been set. In addition, the slave I/O thread generates a warning if either of the following is true:

- No master having the expected `server_uuid` exists.
- The master's `server_uuid` has changed, although no `CHANGE MASTER TO` statement has ever been executed.

18.1.6.1 Replication and Binary Logging Option and Variable Reference

The following tables list basic information about the MySQL command-line options and system variables applicable to replication and the binary log.

Table 18.3 Summary of Replication options and variables in MySQL 8.0

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
`abort-slave-event-count`		
Yes	No	No
Yes		No
DESCRIPTION: Option used by mysql-test for debugging and testing of replication		
`binlog_gtid_simple_recovery`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Controls how binary logs are iterated during GTID recovery		
`Com_change_master`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of CHANGE MASTER TO statements		
`Com_show_master_status`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW MASTER STATUS statements		
`Com_show_new_master`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW NEW MASTER statements		
`Com_show_slave_hosts`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW SLAVE HOSTS statements		
`Com_show_slave_status`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW SLAVE STATUS statements		
`Com_show_slave_status_nonblocking`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW SLAVE STATUS NONBLOCKING statements		
`Com_slave_start`		
No	No	Yes

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
No	Both	No
DESCRIPTION: Count of START SLAVE statements		
`Com_slave_stop`		
No	No	Yes
No	Both	No
DESCRIPTION: Count of STOP SLAVE statements		
`disconnect-slave-event-count`		
Yes	No	No
Yes		No
DESCRIPTION: Option used by mysql-test for debugging and testing of replication		
`enforce-gtid-consistency`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Prevents execution of statements that cannot be logged in a transactionally safe manner		
`enforce_gtid_consistency`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Prevents execution of statements that cannot be logged in a transactionally safe manner		
`executed-gtids-compression-period`		
Yes	No	No
Yes		No
DESCRIPTION: Deprecated and will be removed in a future version. Use the renamed gtid-executed-compression-period instead.		
`executed_gtids_compression_period`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Deprecated and will be removed in a future version. Use the renamed gtid_executed_compression_period instead.		
`gtid-executed-compression-period`		
Yes	No	No
Yes		No
DESCRIPTION: Compress gtid_executed table each time this many transactions have occurred. 0 means never compress this table. Applies only when binary logging is disabled.		
`gtid-mode`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Controls whether GTID based logging is enabled and what type of transactions the logs can contain		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
gtid_executed		
No	Yes	No
No	Global	No
DESCRIPTION: Global: All GTIDs in the binary log (global) or current transaction (session). Read-only.		
gtid_executed_compression_period		
No	Yes	No
No	Global	Yes
DESCRIPTION: Compress gtid_executed table each time this many transactions have occurred. 0 means never compress this table. Applies only when binary logging is disabled.		
gtid_mode		
No	Yes	No
No	Global	Yes
DESCRIPTION: Controls whether GTID based logging is enabled and what type of transactions the logs can contain		
gtid_next		
No	Yes	No
No	Session	Yes
DESCRIPTION: Specifies the GTID for the next statement to execute. See documentation for details.		
gtid_owned		
No	Yes	No
No	Both	No
DESCRIPTION: The set of GTIDs owned by this client (session), or by all clients, together with the thread ID of the owner (global). Read-only.		
gtid_purged		
No	Yes	No
No	Global	Yes
DESCRIPTION: The set of all GTIDs that have been purged from the binary log.		
init_slave		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Statements that are executed when a slave connects to a master		
log-slave-updates		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Tells the slave to log the updates performed by its SQL thread to its own binary log		
log_slave_updates		
Yes	Yes	No
Yes	Global	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Whether the slave should log the updates performed by its SQL thread to its own binary log. Read-only; set using the --log-slave-updates server option.		
log_statements_unsafe_for_binlog		
No	Yes	No
No	Global	Yes
DESCRIPTION: Disables error 1592 warnings being written to the error log		
master-info-file		
Yes	No	No
Yes		No
DESCRIPTION: The location and name of the file that remembers the master and where the I/O replication thread is in the master's binary logs		
master-info-repository		
Yes	No	No
Yes		No
DESCRIPTION: Whether to write master status information and replication I/O thread location in the master's binary logs to a file or table.		
master-retry-count		
Yes	No	No
Yes		No
DESCRIPTION: Number of tries the slave makes to connect to the master before giving up		
master_info_repository		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Whether to write master status information and replication I/O thread location in the master's binary logs to a file or table		
relay-log		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: The location and base name to use for relay logs		
relay-log-index		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: The location and name to use for the file that keeps a list of the last relay logs		
relay-log-info-file		
Yes	No	No
Yes		No
DESCRIPTION: The location and name of the file that remembers where the SQL replication thread is in the relay logs		
relay-log-info-repository		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	No	No
Yes		No
DESCRIPTION: Whether to write the replication SQL thread's location in the relay logs to a file or a table		
`relay-log-recovery`		
Yes	No	No
Yes		No
DESCRIPTION: Enables automatic recovery of relay log files from master at startup		
`relay_log_basename`		
No	Yes	No
No	Global	No
DESCRIPTION: Complete path to relay log, including filename		
`relay_log_index`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: The name of the relay log index file		
`relay_log_info_file`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: The name of the file in which the slave records information about the relay logs		
`relay_log_info_repository`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Whether to write the replication SQL thread's location in the relay logs to a file or a table		
`relay_log_purge`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Determines whether relay logs are purged		
`relay_log_recovery`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Whether automatic recovery of relay log files from master at startup is enabled; must be enabled for a crash-safe slave.		
`relay_log_space_limit`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Maximum space to use for all relay logs		
`replicate-do-db`		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave SQL thread to restrict replication to the specified database		
`replicate-do-table`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave SQL thread to restrict replication to the specified table		
`replicate-ignore-db`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave SQL thread not to replicate to the specified database		
`replicate-ignore-table`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave SQL thread not to replicate to the specified table		
`replicate-rewrite-db`		
Yes	No	No
Yes		No
DESCRIPTION: Updates to a database with a different name than the original		
`replicate-same-server-id`		
Yes	No	No
Yes		No
DESCRIPTION: In replication, if set to 1, do not skip events having our server id		
`replicate-wild-do-table`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave thread to restrict replication to the tables that match the specified wildcard pattern		
`replicate-wild-ignore-table`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the slave thread not to replicate to the tables that match the given wildcard pattern		
`report-host`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Host name or IP of the slave to be reported to the master during slave registration		
`report-password`		
Yes	Yes	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	Global	No
DESCRIPTION: An arbitrary password that the slave server should report to the master. Not the same as the password for the MySQL replication user account.		
`report-port`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Port for connecting to slave reported to the master during slave registration		
`report-user`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: An arbitrary user name that a slave server should report to the master. Not the same as the name used with the MySQL replication user account.		
`Rpl_semi_sync_master_clients`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of semisynchronous slaves		
`rpl_semi_sync_master_enabled`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Whether semisynchronous replication is enabled on the master		
`Rpl_semi_sync_master_net_avg_wait_time`		
No	No	Yes
No	Global	No
DESCRIPTION: The average time the master waited for a slave reply		
`Rpl_semi_sync_master_net_wait_time`		
No	No	Yes
No	Global	No
DESCRIPTION: The total time the master waited for slave replies		
`Rpl_semi_sync_master_net_waits`		
No	No	Yes
No	Global	No
DESCRIPTION: The total number of times the master waited for slave replies		
`Rpl_semi_sync_master_no_times`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of times the master turned off semisynchronous replication		
`Rpl_semi_sync_master_no_tx`		
No	No	Yes
No	Global	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Number of commits not acknowledged successfully `Rpl_semi_sync_master_status`		
No	No	Yes
No	Global	No
DESCRIPTION: Whether semisynchronous replication is operational on the master `Rpl_semi_sync_master_timefunc_failures`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of times the master failed when calling time functions `rpl_semi_sync_master_timeout`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Number of milliseconds to wait for slave acknowledgment `rpl_semi_sync_master_trace_level`		
No	Yes	No
No	Global	Yes
DESCRIPTION: The semisynchronous replication debug trace level on the master `Rpl_semi_sync_master_tx_avg_wait_time`		
No	No	Yes
No	Global	No
DESCRIPTION: The average time the master waited for each transaction `Rpl_semi_sync_master_tx_wait_time`		
No	No	Yes
No	Global	No
DESCRIPTION: The total time the master waited for transactions `Rpl_semi_sync_master_tx_waits`		
No	No	Yes
No	Global	No
DESCRIPTION: The total number of times the master waited for transactions `rpl_semi_sync_master_wait_for_slave_count`		
No	Yes	No
No	Global	Yes
DESCRIPTION: How many slave acknowledgments the master must receive per transaction before proceeding `rpl_semi_sync_master_wait_no_slave`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Whether master waits for timeout even with no slaves		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
`rpl_semi_sync_master_wait_point`		
No	Yes	No
No	Global	Yes
DESCRIPTION: The wait point for slave transaction receipt acknowledgment		
`Rpl_semi_sync_master_wait_pos_backtraverse`		
No	No	Yes
No	Global	No
DESCRIPTION: The total number of times the master waited for an event with binary coordinates lower than events waited for previously		
`Rpl_semi_sync_master_wait_sessions`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of sessions currently waiting for slave replies		
`Rpl_semi_sync_master_yes_tx`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of commits acknowledged successfully		
`rpl_semi_sync_slave_enabled`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Whether semisynchronous replication is enabled on slave		
`Rpl_semi_sync_slave_status`		
No	No	Yes
No	Global	No
DESCRIPTION: Whether semisynchronous replication is operational on slave		
`rpl_semi_sync_slave_trace_level`		
No	Yes	No
No	Global	Yes
DESCRIPTION: The semisynchronous replication debug trace level on the slave		
`rpl_stop_slave_timeout`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Set the number of seconds that STOP SLAVE waits before timing out.		
`server_uuid`		
No	Yes	No
No	Global	No
DESCRIPTION: The server's globally unique ID, automatically (re)generated at server start		
`show-slave-auth-info`		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	No	No
Yes		No
DESCRIPTION: Show user name and password in SHOW SLAVE HOSTS on this master		
`simplified_binlog_gtid_recovery`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Controls how binary logs are iterated during GTID recovery		
`skip-slave-start`		
Yes	No	No
Yes		No
DESCRIPTION: If set, slave is not autostarted		
`slave-checkpoint-group`		
Yes	No	No
Yes		No
DESCRIPTION: Maximum number of transactions processed by a multi-threaded slave before a checkpoint operation is called to update progress status. Not supported by MySQL Cluster.		
`slave-checkpoint-period`		
Yes	No	No
Yes		No
DESCRIPTION: Update progress status of multi-threaded slave and flush relay log info to disk after this number of milliseconds. Not supported by MySQL Cluster.		
`slave-load-tmpdir`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: The location where the slave should put its temporary files when replicating a LOAD DATA INFILE statement		
`slave-max-allowed-packet`		
Yes	No	No
Yes		No
DESCRIPTION: Maximum size, in bytes, of a packet that can be sent from a replication master to a slave; overrides max_allowed_packet.		
`slave_net_timeout`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Number of seconds to wait for more data from a master/slave connection before aborting the read		
`slave-parallel-type`		
Yes	No	No
Yes		No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Tells the slave to use database partioning (DATABASE) or timestamp information (LOGICAL_CLOCK) from the master to parallelize transactions. The default is DATABASE.		
`slave-parallel-workers`		
Yes	No	No
Yes		No
DESCRIPTION: Number of worker threads for executing events in parallel. Set to 0 (the default) to disable slave multi-threading. Not supported by MySQL Cluster.		
`slave-pending-jobs-size-max`		
Yes	No	No
No		No
DESCRIPTION: Maximum size of slave worker queues holding events not yet applied.		
`slave-rows-search-algorithms`		
Yes	No	No
Yes		No
DESCRIPTION: Determines search algorithms used for slave update batching. Any 2 or 3 from the list INDEX_SEARCH, TABLE_SCAN, HASH_SCAN; the default is TABLE_SCAN,INDEX_SCAN.		
`slave-skip-errors`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Tells the slave thread to continue replication when a query returns an error from the provided list		
`slave_checkpoint_group`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Maximum number of transactions processed by a multi-threaded slave before a checkpoint operation is called to update progress status. Not supported by MySQL Cluster.		
`slave_checkpoint_period`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Update progress status of multi-threaded slave and flush relay log info to disk after this number of milliseconds. Not supported by MySQL Cluster.		
`slave_compressed_protocol`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Use compression on master/slave protocol		
`slave_exec_mode`		
Yes	Yes	No
Yes	Global	Yes

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Allows for switching the slave thread between IDEMPOTENT mode (key and some other errors suppressed) and STRICT mode; STRICT mode is the default, except for MySQL Cluster, where IDEMPOTENT is always used		
`Slave_heartbeat_period`		
No	No	Yes
No	Global	No
DESCRIPTION: The slave's replication heartbeat interval, in seconds		
`slave_max_allowed_packet`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Maximum size, in bytes, of a packet that can be sent from a replication master to a slave; overrides max_allowed_packet.		
`Slave_open_temp_tables`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of temporary tables that the slave SQL thread currently has open		
`slave_parallel_type`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Tells the slave to use database partioning (DATABASE) or information (LOGICAL_CLOCK) from master to parallelize transactions. The default is DATABASE.		
`slave_parallel_workers`		
Yes	Yes	No
No	Global	Yes
DESCRIPTION: Number of worker threads for executing events in parallel. Set to 0 (the default) to disable slave multi-threading. Not supported by MySQL Cluster.		
`slave_pending_jobs_size_max`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Maximum size of slave worker queues holding events not yet applied.		
`slave_preserve_commit_order`		
Yes	Yes	No
No	Global	Yes
DESCRIPTION: Ensures that all commits by slave workers happen in the same order as on the master to maintain consistency when using parallel worker threads.		
`Slave_retried_transactions`		
No	No	Yes
No	Global	No
DESCRIPTION: The total number of times since startup that the replication slave SQL thread has retried transactions		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
slave_rows_search_algorithms		
No	Yes	No
No	Global	Yes
DESCRIPTION: Determines search algorithms used for slave update batching. Any 2 or 3 from the list INDEX_SEARCH, TABLE_SCAN, HASH_SCAN; the default is TABLE_SCAN,INDEX_SCAN.		
Slave_running		
No	No	Yes
No	Global	No
DESCRIPTION: The state of this server as a replication slave (slave I/O thread status)		
slave_transaction_retries		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Number of times the slave SQL thread will retry a transaction in case it failed with a deadlock or elapsed lock wait timeout, before giving up and stopping		
slave_type_conversions		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Controls type conversion mode on replication slave. Value is a list of zero or more elements from the list: ALL_LOSSY, ALL_NON_LOSSY. Set to an empty string to disallow type conversions between master and slave.		
sql_slave_skip_counter		
No	Yes	No
No	Global	Yes
DESCRIPTION: Number of events from the master that a slave server should skip. Not compatible with GTID replication.		
sync_binlog		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Synchronously flush binary log to disk after every #th event		
sync_master_info		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Synchronize master.info to disk after every #th event.		
sync_relay_log		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Synchronize relay log to disk after every #th event.		
sync_relay_log_info		
Yes	Yes	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	Global	Yes
DESCRIPTION: Synchronize relay.info file to disk after every #th event.		

Section 18.1.6.2, “Replication Master Options and Variables”, provides more detailed information about options and variables relating to replication master servers. For more information about options and variables relating to replication slaves, see Section 18.1.6.3, “Replication Slave Options and Variables”.

Table 18.4 Summary of Binary Logging options and variables in MySQL 8.0

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
`binlog-checksum`		
Yes	No	No
Yes		No
DESCRIPTION: Enable/disable binary log checksums		
`binlog-do-db`		
Yes	No	No
Yes		No
DESCRIPTION: Limits binary logging to specific databases		
`binlog_format`		
Yes	Yes	No
Yes	Both	Yes
DESCRIPTION: Specifies the format of the binary log		
`binlog-ignore-db`		
Yes	No	No
Yes		No
DESCRIPTION: Tells the master that updates to the given database should not be logged to the binary log		
`binlog-row-event-max-size`		
Yes	No	No
Yes		No
DESCRIPTION: Binary log max event size		
`binlog-rows-query-log-events`		
Yes	No	No
Yes		No
DESCRIPTION: Enables logging of rows query log events when using row-based logging. Disabled by default. Do not enable when producing logs for pre-5.6.2 slaves/readers.		
`Binlog_cache_disk_use`		
No	No	Yes
No	Global	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Number of transactions that used a temporary file instead of the binary log cache `binlog_cache_size`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Size of the cache to hold the SQL statements for the binary log during a transaction `Binlog_cache_use`		
No	No	Yes
No	Global	No
DESCRIPTION: Number of transactions that used the temporary binary log cache `binlog_checksum`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Enable/disable binary log checksums `binlog_direct_non_transactional_updates`		
Yes	Yes	No
Yes	Both	Yes
DESCRIPTION: Causes updates using statement format to nontransactional engines to be written directly to binary log. See documentation before using. `binlog_error_action`		
Yes	Yes	No
Yes	Both	Yes
DESCRIPTION: Controls what happens when the server cannot write to the binary log. `binlog_group_commit_sync_delay`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Sets the number of microseconds to wait before synchronizing transactions to disk. `binlog_group_commit_sync_no_delay_count`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Sets the maximum number of transactions to wait for before aborting the current delay specified by binlog_group_commit_sync_delay. `binlog_max_flush_queue_time`		
No	Yes	No
No	Global	Yes
DESCRIPTION: How long to read transactions before flushing to binary log `binlog_order_commits`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Whether to commit in same order as writes to binary log		

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
binlog_row_image		
Yes	Yes	No
Yes	Both	Yes
DESCRIPTION: Use full or minimal images when logging row changes. Allowed values are full, minimal, and noblob.		
binlog_rows_query_log_events		
No	Yes	No
No	Both	Yes
DESCRIPTION: When TRUE, enables logging of rows query log events in row-based logging mode. FALSE by default. Do not enable when producing logs for pre-5.6.2 replication slaves or other readers.		
Binlog_stmt_cache_disk_use		
No	No	Yes
No	Global	No
DESCRIPTION: Number of nontransactional statements that used a temporary file instead of the binary log statement cache		
binlog_stmt_cache_size		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Size of the cache to hold nontransactional statements for the binary log during a transaction		
Binlog_stmt_cache_use		
No	No	Yes
No	Global	No
DESCRIPTION: Number of statements that used the temporary binary log statement cache		
binlogging_impossible_mode		
Yes	Yes	No
Yes	Both	Yes
DESCRIPTION: Deprecated and will be removed in a future version. Use the renamed binlog_error_action instead.		
Com_show_binlog_events		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW BINLOG EVENTS statements		
Com_show_binlogs		
No	No	Yes
No	Both	No
DESCRIPTION: Count of SHOW BINLOGS statements		
log-bin-use-v1-row-events		
Yes	Yes	No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
Yes	Global	No
DESCRIPTION: Use version 1 binary log row events		
`log_bin_basename`		
No	Yes	No
No	Global	No
DESCRIPTION: Complete path to binary log, including filename		
`log_bin_use_v1_row_events`		
Yes	Yes	No
Yes	Global	No
DESCRIPTION: Shows whether server is using version 1 binary log row events		
`master-verify-checksum`		
Yes	No	No
Yes		No
DESCRIPTION: Cause master to examine checksums when reading from the binary log		
`master_verify_checksum`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Cause master to read checksums from binary log.		
`max-binlog-dump-events`		
Yes	No	No
Yes		No
DESCRIPTION: Option used by mysql-test for debugging and testing of replication		
`max_binlog_cache_size`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Can be used to restrict the total size used to cache a multi-statement transaction		
`max_binlog_size`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Binary log will be rotated automatically when size exceeds this value		
`max_binlog_stmt_cache_size`		
Yes	Yes	No
Yes	Global	Yes
DESCRIPTION: Can be used to restrict the total size used to cache all nontransactional statements during a transaction		
`slave-sql-verify-checksum`		
Yes	No	No
Yes		No

Option or Variable Name		
Command Line	**System Variable**	**Status Variable**
Option File	**Scope**	**Dynamic**
Notes		
DESCRIPTION: Cause slave to examine checksums when reading from the relay log		
`slave_sql_verify_checksum`		
No	Yes	No
No	Global	Yes
DESCRIPTION: Cause slave to examine checksums when reading from relay log.		
`sporadic-binlog-dump-fail`		
Yes	No	No
Yes		No
DESCRIPTION: Option used by mysql-test for debugging and testing of replication		

Section 18.1.6.4, “Binary Logging Options and Variables”, provides more detailed information about options and variables relating to binary logging. For additional general information about the binary log, see Section 5.4.4, “The Binary Log”.

For information about the `sql_log_bin` and `sql_log_off` variables, see Section 5.1.5, “Server System Variables”.

For a table showing *all* command-line options, system and status variables used with `mysqld`, see Section 5.1.3, “Server Option and Variable Reference”.

18.1.6.2 Replication Master Options and Variables

This section describes the server options and system variables that you can use on replication master servers. You can specify the options either on the command line or in an option file. You can specify system variable values using `SET`.

On the master and each slave, you must use the `server-id` option to establish a unique replication ID. For each server, you should pick a unique positive integer in the range from 1 to $2^{32} - 1$, and each ID must be different from every other ID in use by any other replication master or slave. Example: `server-id=3`.

For options used on the master for controlling binary logging, see Section 18.1.6.4, “Binary Logging Options and Variables”.

System Variables Used on Replication Masters

The following system variables are used to control replication masters:

- `auto_increment_increment`

System Variable	**Name**	`auto_increment_increment`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1`
	Min Value	`1`

	Max Value	65535

`auto_increment_increment` and `auto_increment_offset` are intended for use with master-to-master replication, and can be used to control the operation of `AUTO_INCREMENT` columns. Both variables have global and session values, and each can assume an integer value between 1 and 65,535 inclusive. Setting the value of either of these two variables to 0 causes its value to be set to 1 instead. Attempting to set the value of either of these two variables to an integer greater than 65,535 or less than 0 causes its value to be set to 65,535 instead. Attempting to set the value of `auto_increment_increment` or `auto_increment_offset` to a noninteger value produces an error, and the actual value of the variable remains unchanged.

Note

`auto_increment_increment` is intended for use with MySQL Cluster, which is not currently supported in MySQL 8.0.

These two variables affect `AUTO_INCREMENT` column behavior as follows:

- `auto_increment_increment` controls the interval between successive column values. For example:

```
mysql> SHOW VARIABLES LIKE 'auto_inc%';
+--------------------------+-------+
| Variable_name            | Value |
+--------------------------+-------+
| auto_increment_increment | 1     |
| auto_increment_offset    | 1     |
+--------------------------+-------+
2 rows in set (0.00 sec)

mysql> CREATE TABLE autoinc1
    -> (col INT NOT NULL AUTO_INCREMENT PRIMARY KEY);
  Query OK, 0 rows affected (0.04 sec)

mysql> SET @@auto_increment_increment=10;
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW VARIABLES LIKE 'auto_inc%';
+--------------------------+-------+
| Variable_name            | Value |
+--------------------------+-------+
| auto_increment_increment | 10    |
| auto_increment_offset    | 1     |
+--------------------------+-------+
2 rows in set (0.01 sec)

mysql> INSERT INTO autoinc1 VALUES (NULL), (NULL), (NULL), (NULL);
Query OK, 4 rows affected (0.00 sec)
Records: 4  Duplicates: 0  Warnings: 0

mysql> SELECT col FROM autoinc1;
+-----+
| col |
+-----+
|   1 |
|  11 |
|  21 |
|  31 |
+-----+
4 rows in set (0.00 sec)
```

- `auto_increment_offset` determines the starting point for the `AUTO_INCREMENT` column value. Consider the following, assuming that these statements are executed during the same session as the example given in the description for `auto_increment_increment`:

```
mysql> SET @@auto_increment_offset=5;
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW VARIABLES LIKE 'auto_inc%';
+--------------------------+-------+
| Variable_name            | Value |
+--------------------------+-------+
| auto_increment_increment | 10    |
| auto_increment_offset    | 5     |
+--------------------------+-------+
2 rows in set (0.00 sec)

mysql> CREATE TABLE autoinc2
    -> (col INT NOT NULL AUTO_INCREMENT PRIMARY KEY);
Query OK, 0 rows affected (0.06 sec)

mysql> INSERT INTO autoinc2 VALUES (NULL), (NULL), (NULL), (NULL);
Query OK, 4 rows affected (0.00 sec)
Records: 4  Duplicates: 0  Warnings: 0

mysql> SELECT col FROM autoinc2;
+-----+
| col |
+-----+
|   5 |
|  15 |
|  25 |
|  35 |
+-----+
4 rows in set (0.02 sec)
```

When the value of `auto_increment_offset` is greater than that of `auto_increment_increment`, the value of `auto_increment_offset` is ignored.

If either of these variables is changed, and then new rows inserted into a table containing an `AUTO_INCREMENT` column, the results may seem counterintuitive because the series of `AUTO_INCREMENT` values is calculated without regard to any values already present in the column, and the next value inserted is the least value in the series that is greater than the maximum existing value in the `AUTO_INCREMENT` column. The series is calculated like this:

`auto_increment_offset` + *N* × `auto_increment_increment`

where *N* is a positive integer value in the series [1, 2, 3, ...]. For example:

```
mysql> SHOW VARIABLES LIKE 'auto_inc%';
+--------------------------+-------+
| Variable_name            | Value |
+--------------------------+-------+
| auto_increment_increment | 10    |
| auto_increment_offset    | 5     |
+--------------------------+-------+
2 rows in set (0.00 sec)

mysql> SELECT col FROM autoinc1;
+-----+
| col |
+-----+
|   1 |
|  11 |
|  21 |
|  31 |
+-----+
4 rows in set (0.00 sec)

mysql> INSERT INTO autoinc1 VALUES (NULL), (NULL), (NULL), (NULL);
Query OK, 4 rows affected (0.00 sec)
Records: 4  Duplicates: 0  Warnings: 0
```

```
mysql> SELECT col FROM autoinc1;
+-----+
| col |
+-----+
|   1 |
|  11 |
|  21 |
|  31 |
|  35 |
|  45 |
|  55 |
|  65 |
+-----+
8 rows in set (0.00 sec)
```

The values shown for `auto_increment_increment` and `auto_increment_offset` generate the series 5 + *N* × 10, that is, [5, 15, 25, 35, 45, ...]. The highest value present in the `col` column prior to the `INSERT` is 31, and the next available value in the `AUTO_INCREMENT` series is 35, so the inserted values for `col` begin at that point and the results are as shown for the `SELECT` query.

It is not possible to restrict the effects of these two variables to a single table; these variables control the behavior of all `AUTO_INCREMENT` columns in *all* tables on the MySQL server. If the global value of either variable is set, its effects persist until the global value is changed or overridden by setting the session value, or until `mysqld` is restarted. If the local value is set, the new value affects `AUTO_INCREMENT` columns for all tables into which new rows are inserted by the current user for the duration of the session, unless the values are changed during that session.

The default value of `auto_increment_increment` is 1. See Section 18.4.1.1, “Replication and AUTO_INCREMENT”.

- `auto_increment_offset`

System Variable	Name	`auto_increment_offset`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`1`
	Min Value	`1`
	Max Value	`65535`

This variable has a default value of 1. For more information, see the description for `auto_increment_increment`.

Note

`auto_increment_offset` is intended for use with MySQL Cluster, which is not currently supported in MySQL 8.0.

18.1.6.3 Replication Slave Options and Variables

This section explains the server options and system variables that apply to slave replication servers and contains the following:

Startup Options for Replication Slaves

Options for Logging Slave Status to Tables

System Variables Used on Replication Slaves

Specify the options either on the command line or in an option file. Many of the options can be set while the server is running by using the `CHANGE MASTER TO` statement. Specify system variable values using `SET`.

Server ID. On the master and each slave, you must use the `server-id` option to establish a unique replication ID in the range from 1 to $2^{32} - 1$. "Unique" means that each ID must be different from every other ID in use by any other replication master or slave. Example `my.cnf` file:

```
[mysqld]
server-id=3
```

Startup Options for Replication Slaves

This section explains startup options for controlling replication slave servers. Many of these options can be set while the server is running by using the `CHANGE MASTER TO` statement. Others, such as the `--replicate-*` options, can be set only when the slave server starts. Replication-related system variables are discussed later in this section.

- `--log-slave-updates`

Command-Line Format	`--log-slave-updates`	
System Variable	**Name**	`log_slave_updates`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`OFF`

Normally, a slave does not write any updates that are received from a master server to its own binary log. This option causes the slave to write the updates performed by its SQL thread to its own binary log. For this option to have any effect, the slave must also be started with the `--log-bin` option to enable binary logging. `--log-slave-updates` is used when you want to chain replication servers. For example, you might want to set up replication servers using this arrangement:

```
A -> B -> C
```

Here, `A` serves as the master for the slave `B`, and `B` serves as the master for the slave `C`. For this to work, `B` must be both a master *and* a slave. You must start both `A` and `B` with `--log-bin` to enable binary logging, and `B` with the `--log-slave-updates` option so that updates received from `A` are logged by `B` to its binary log.

- `--log-warnings[=level]`

Deprecated	5.7.2	
Command-Line Format	`--log-warnings[=#]`	
System Variable	**Name**	`log_warnings`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer

	Default	2
	Min Value	0
	Max Value	4294967295
Permitted Values (64-bit platforms)	**Type**	integer
	Default	2
	Min Value	0
	Max Value	18446744073709551615

Note

The `log_error_verbosity` system variable is preferred over, and should be used instead of, the `--log-warnings` option or `log_warnings` system variable. For more information, see the descriptions of `log_error_verbosity` and `log_warnings`. The `--log-warnings` command-line option and `log_warnings` system variable are deprecated and will be removed in a future MySQL release.

Causes the server to record more messages to the error log about what it is doing. With respect to replication, the server generates warnings that it succeeded in reconnecting after a network or connection failure, and provides information about how each slave thread started. This variable is enabled by default with a value of 2. To disable it, set it to 0. The server logs messages about statements that are unsafe for statement-based logging if the value is greater than 0. Aborted connections and access-denied errors for new connection attempts are logged if the value is greater than 1. See Section B.5.2.10, “Communication Errors and Aborted Connections”.

Note

The effects of this option are not limited to replication. It produces warnings across a spectrum of server activities.

- `--master-info-file=`*`file_name`*

Command-Line Format	`--master-info-file=file_name`	
Permitted Values	**Type**	file name
	Default	`master.info`

The name to use for the file in which the slave records information about the master. The default name is `master.info` in the data directory. For information about the format of this file, see Section 18.2.4.2, “Slave Status Logs”.

- `--master-retry-count=`*`count`*

Deprecated	5.6.1	
Command-Line Format	`--master-retry-count=#`	
Permitted Values (32-bit platforms)	**Type**	integer
	Default	86400
	Min Value	0
	Max Value	4294967295

Permitted Values (64-bit platforms)	**Type**	integer
	Default	`86400`
	Min Value	`0`
	Max Value	`18446744073709551615`

The number of times that the slave tries to connect to the master before giving up. Reconnects are attempted at intervals set by the `MASTER_CONNECT_RETRY` option of the `CHANGE MASTER TO` statement (default 60). Reconnects are triggered when data reads by the slave time out according to the `--slave-net-timeout` option. The default value is 86400. A value of 0 means “infinite”; the slave attempts to connect forever.

This option is deprecated and will be removed in a future MySQL release. Applications should be updated to use the `MASTER_RETRY_COUNT` option of the `CHANGE MASTER TO` statement instead.

- `--max-relay-log-size=`*`size`*

Command-Line Format	`--max-relay-log-size=#`	
System Variable	**Name**	`max_relay_log_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1073741824`

The size at which the server rotates relay log files automatically. If this value is nonzero, the relay log is rotated automatically when its size exceeds this value. If this value is zero (the default), the size at which relay log rotation occurs is determined by the value of `max_binlog_size`. For more information, see Section 18.2.4.1, “The Slave Relay Log”.

- `--relay-log=`*`file_name`*

Command-Line Format	`--relay-log=file_name`	
System Variable	**Name**	`relay_log`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The base name for the relay log. For the default replication channel, the default base name for relay logs is *`host_name`*`-relay-bin`. For non-default replication channels, the default base name for relay logs is *`host_name-channel`*`-relay-bin`, where *`channel`* is the name of the replication channel recorded in this relay log. The server writes the file in the data directory unless the base name is given with a leading absolute path name to specify a different directory. The server creates relay log files in sequence by adding a numeric suffix to the base name.

Due to the manner in which MySQL parses server options, if you specify this option, you must supply a value; *the default base name is used only if the option is not actually specified*. If you use the `--relay-log` option without specifying a value, unexpected behavior is likely to result; this behavior depends on the other options used, the order in which they are specified, and whether they are specified on the command line or in an option file. For more information about how MySQL handles server options, see Section 4.2.3, “Specifying Program Options”.

If you specify this option, the value specified is also used as the base name for the relay log index file. You can override this behavior by specifying a different relay log index file base name using the `--relay-log-index` option.

When the server reads an entry from the index file, it checks whether the entry contains a relative path. If it does, the relative part of the path is replaced with the absolute path set using the `--relay-log` option. An absolute path remains unchanged; in such a case, the index must be edited manually to enable the new path or paths to be used. Previously, manual intervention was required whenever relocating the binary log or relay log files. (Bug #11745230, Bug #12133)

You may find the `--relay-log` option useful in performing the following tasks:

- Creating relay logs whose names are independent of host names.
- If you need to put the relay logs in some area other than the data directory because your relay logs tend to be very large and you do not want to decrease `max_relay_log_size`.
- To increase speed by using load-balancing between disks.

You can obtain the relay log file name (and path) from the `relay_log_basename` system variable.

- `--relay-log-index=`*`file_name`*

Command-Line Format	`--relay-log-index=file_name`	
System Variable	**Name**	`relay_log_index`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The name to use for the relay log index file. The default name is *`host_name`*`-relay-bin.index` in the data directory, where *`host_name`* is the name of the server. For the default replication channel, the default name is *`host_name`*`-relay-bin.index`. For non-default replication channels, the default name is *`host_name`*`-`*`channel`*`-relay-bin.index`, where *`channel`* is the name of the replication channel recorded in this relay log index.

Due to the manner in which MySQL parses server options, if you specify this option, you must supply a value; *the default base name is used only if the option is not actually specified*. If you use the `--relay-log-index` option without specifying a value, unexpected behavior is likely to result; this behavior depends on the other options used, the order in which they are specified, and whether they are specified on the command line or in an option file. For more information about how MySQL handles server options, see Section 4.2.3, “Specifying Program Options”.

If you specify this option, the value specified is also used as the base name for the relay logs. You can override this behavior by specifying a different relay log file base name using the `--relay-log` option.

- `--relay-log-info-file=`*`file_name`*

Command-Line Format	--relay-log-info-file=file_name	
Permitted Values	Type	file name
	Default	relay-log.info

The name to use for the file in which the slave records information about the relay logs. The default name is relay-log.info in the data directory. For information about the format of this file, see Section 18.2.4.2, "Slave Status Logs".

- --relay-log-purge={0|1}

Command-Line Format	--relay-log-purge	
System Variable	Name	relay_log_purge
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	TRUE

Disable or enable automatic purging of relay logs as soon as they are no longer needed. The default value is 1 (enabled). This is a global variable that can be changed dynamically with SET GLOBAL relay_log_purge = *N*. Disabling purging of relay logs when using the --relay-log-recovery option risks data consistency and is therefore not crash-safe.

- --relay-log-recovery={0|1}

Command-Line Format	--relay-log-recovery	
Permitted Values	Type	boolean
	Default	FALSE

Enables automatic relay log recovery immediately following server startup. The recovery process creates a new relay log file, initializes the SQL thread position to this new relay log, and initializes the I/O thread to the SQL thread position. Reading of the relay log from the master then continues. This should be used following a crash on the replication slave to ensure that no possibly corrupted relay logs are processed. The default value is 0 (disabled).

To provide a crash-safe slave, this option must be enabled (set to 1), --relay-log-info-repository must be set to TABLE, and relay-log-purge must be enabled. Enabling the --relay-log-recovery option when relay-log-purge is disabled risks reading the relay log from files that were not purged, leading to data inconsistency, and is therefore not crash-safe. See Crash-safe replication, for more information.

When using a multi-threaded slave (in other words slave_parallel_workers is greater than 0), inconsistencies such as gaps can occur in the sequence of transactions that have been executed from the relay log. Enabling the --relay-log-recovery option when there are inconsistencies causes an error and the option has no effect. The solution in this situation is to issue START SLAVE UNTIL SQL_AFTER_MTS_GAPS, which brings the server to a more consistent state, then issue RESET SLAVE to remove the relay logs. See Section 18.4.1.34, "Replication and Transaction Inconsistencies" for more information.

- --relay-log-space-limit=*size*

Command-Line Format	--relay-log-space-limit=#	
System Variable	Name	relay_log_space_limit

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709551615`

This option places an upper limit on the total size in bytes of all relay logs on the slave. A value of 0 means “no limit”. This is useful for a slave server host that has limited disk space. When the limit is reached, the I/O thread stops reading binary log events from the master server until the SQL thread has caught up and deleted some unused relay logs. Note that this limit is not absolute: There are cases where the SQL thread needs more events before it can delete relay logs. In that case, the I/O thread exceeds the limit until it becomes possible for the SQL thread to delete some relay logs because not doing so would cause a deadlock. You should not set `--relay-log-space-limit` to less than twice the value of `--max-relay-log-size` (or `--max-binlog-size` if `--max-relay-log-size` is 0). In that case, there is a chance that the I/O thread waits for free space because `--relay-log-space-limit` is exceeded, but the SQL thread has no relay log to purge and is unable to satisfy the I/O thread. This forces the I/O thread to ignore `--relay-log-space-limit` temporarily.

- `--replicate-do-db=`*`db_name`*

Command-Line Format	`--replicate-do-db=name`	
Permitted Values	**Type**	string

Creates a replication filter that selects using the name of a database. Such filters can also be created using `CHANGE REPLICATION FILTER REPLICATE_DO_DB`. Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named *`channel_1`* use `--replicate-do-db:`*`channel_1`*`:`*`db_name`*. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, “Replication Channel Based Filters” for more information. The precise effect of this filter depends on whether statement-based or row-based replication is in use.

Statement-based replication. Tell the slave SQL thread to restrict replication to statements where the default database (that is, the one selected by `USE`) is *`db_name`*. To specify more than one database, use this option multiple times, once for each database; however, doing so does *not* replicate cross-database statements such as `UPDATE` *`some_db.some_table`* `SET foo='bar'` while a different database (or no database) is selected.

Warning

To specify multiple databases you *must* use multiple instances of this option. Because database names can contain commas, if you supply a comma separated list then the list is treated as the name of a single database.

An example of what does not work as you might expect when using statement-based replication: If the slave is started with `--replicate-do-db=sales` and you issue the following statements on the master, the `UPDATE` statement is *not* replicated:

```
USE prices;
UPDATE sales.january SET amount=amount+1000;
```

The main reason for this “check just the default database” behavior is that it is difficult from the statement alone to know whether it should be replicated (for example, if you are using multiple-table `DELETE` statements or multiple-table `UPDATE` statements that act across multiple databases). It is also faster to check only the default database rather than all databases if there is no need.

Row-based replication. Tells the slave SQL thread to restrict replication to database *`db_name`*. Only tables belonging to *`db_name`* are changed; the current database has no effect on this. Suppose that the slave is started with `--replicate-do-db=sales` and row-based replication is in effect, and then the following statements are run on the master:

```
USE prices;
UPDATE sales.february SET amount=amount+100;
```

The `february` table in the `sales` database on the slave is changed in accordance with the `UPDATE` statement; this occurs whether or not the `USE` statement was issued. However, issuing the following statements on the master has no effect on the slave when using row-based replication and `--replicate-do-db=sales`:

```
USE prices;
UPDATE prices.march SET amount=amount-25;
```

Even if the statement `USE prices` were changed to `USE sales`, the `UPDATE` statement's effects would still not be replicated.

Another important difference in how `--replicate-do-db` is handled in statement-based replication as opposed to row-based replication occurs with regard to statements that refer to multiple databases. Suppose that the slave is started with `--replicate-do-db=db1`, and the following statements are executed on the master:

```
USE db1;
UPDATE db1.table1 SET col1 = 10, db2.table2 SET col2 = 20;
```

If you are using statement-based replication, then both tables are updated on the slave. However, when using row-based replication, only `table1` is affected on the slave; since `table2` is in a different database, `table2` on the slave is not changed by the `UPDATE`. Now suppose that, instead of the `USE db1` statement, a `USE db4` statement had been used:

```
USE db4;
UPDATE db1.table1 SET col1 = 10, db2.table2 SET col2 = 20;
```

In this case, the `UPDATE` statement would have no effect on the slave when using statement-based replication. However, if you are using row-based replication, the `UPDATE` would change `table1` on the slave, but not `table2`—in other words, only tables in the database named by `--replicate-do-db` are changed, and the choice of default database has no effect on this behavior.

If you need cross-database updates to work, use `--replicate-wild-do-table=db_name.%` instead. See Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”.

Note

This option affects replication in the same manner that `--binlog-do-db` affects binary logging, and the effects of the replication format on how `--replicate-do-db` affects replication behavior are the same as those of the logging format on the behavior of `--binlog-do-db`.

This option has no effect on `BEGIN`, `COMMIT`, or `ROLLBACK` statements.

- `--replicate-ignore-db=db_name`

Command-Line Format	`--replicate-ignore-db=name`	
Permitted Values	Type	string

Creates a replication filter using the name of a database. Such filters can also be created using `CHANGE REPLICATION FILTER REPLICATE_IGNORE_DB`. Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named *`channel_1`* use `--replicate-ignore-db:channel_1:db_name`. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, “Replication Channel Based Filters” for more information.

As with `--replicate-do-db`, the precise effect of this filtering depends on whether statement-based or row-based replication is in use, and are described in the next several paragraphs.

Statement-based replication. Tells the slave SQL thread not to replicate any statement where the default database (that is, the one selected by `USE`) is *`db_name`*.

Row-based replication. Tells the slave SQL thread not to update any tables in the database *`db_name`*. The default database has no effect.

When using statement-based replication, the following example does not work as you might expect. Suppose that the slave is started with `--replicate-ignore-db=sales` and you issue the following statements on the master:

```
USE prices;
UPDATE sales.january SET amount=amount+1000;
```

The `UPDATE` statement *is* replicated in such a case because `--replicate-ignore-db` applies only to the default database (determined by the `USE` statement). Because the `sales` database was specified explicitly in the statement, the statement has not been filtered. However, when using row-based replication, the `UPDATE` statement's effects are *not* propagated to the slave, and the slave's copy of the `sales.january` table is unchanged; in this instance, `--replicate-ignore-db=sales` causes *all* changes made to tables in the master's copy of the `sales` database to be ignored by the slave.

To specify more than one database to ignore, use this option multiple times, once for each database. Because database names can contain commas, if you supply a comma separated list then the list will be treated as the name of a single database.

You should not use this option if you are using cross-database updates and you do not want these updates to be replicated. See Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”.

If you need cross-database updates to work, use `--replicate-wild-ignore-table=db_name.%` instead. See Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”.

Note

This option affects replication in the same manner that `--binlog-ignore-db` affects binary logging, and the effects of the replication format on how `--replicate-ignore-db` affects replication behavior are the same as those of the logging format on the behavior of `--binlog-ignore-db`.

This option has no effect on `BEGIN`, `COMMIT`, or `ROLLBACK` statements.

- `--replicate-do-table=db_name.tbl_name`

Command-Line Format	`--replicate-do-table=name`	
Permitted Values	Type	string

Creates a replication filter by telling the slave SQL thread to restrict replication to a given table. To specify more than one table, use this option multiple times, once for each table. This works for both cross-database updates and default database updates, in contrast to `--replicate-do-db`. See Section 18.2.5, "How Servers Evaluate Replication Filtering Rules". You can also create such a filter by issuing a `CHANGE REPLICATION FILTER REPLICATE_DO_TABLE` statement.

Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named `channel_1` use `--replicate-do-table:channel_1:db_name.tbl_name`. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, "Replication Channel Based Filters" for more information.

This option affects only statements that apply to tables. It does not affect statements that apply only to other database objects, such as stored routines. To filter statements operating on stored routines, use one or more of the `--replicate-*-db` options.

- `--replicate-ignore-table=db_name.tbl_name`

Command-Line Format	`--replicate-ignore-table=name`	
Permitted Values	Type	string

Creates a replication filter by telling the slave SQL thread not to replicate any statement that updates the specified table, even if any other tables might be updated by the same statement. To specify more than one table to ignore, use this option multiple times, once for each table. This works for cross-database updates, in contrast to `--replicate-ignore-db`. See Section 18.2.5, "How Servers Evaluate Replication Filtering Rules". You can also create such a filter by issuing a `CHANGE REPLICATION FILTER REPLICATE_IGNORE_TABLE` statement.

Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named `channel_1` use `--replicate-ignore-table:channel_1:db_name.tbl_name`. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, "Replication Channel Based Filters" for more information.

This option affects only statements that apply to tables. It does not affect statements that apply only to other database objects, such as stored routines. To filter statements operating on stored routines, use one or more of the `--replicate-*-db` options.

- `--replicate-rewrite-db=from_name->to_name`

Command-Line Format	`--replicate-rewrite-db=old_name->new_name`	
Permitted Values	Type	string

Tells the slave to create a replication filter that translates the default database (that is, the one selected by `USE`) to *`to_name`* if it was *`from_name`* on the master. Only statements involving tables are affected (not statements such as `CREATE DATABASE`, `DROP DATABASE`, and `ALTER DATABASE`), and only if *`from_name`* is the default database on the master. To specify multiple rewrites, use this option multiple times. The server uses the first one with a *`from_name`* value that matches. The database name translation is done *before* the `--replicate-*` rules are tested. You can also create such a filter by issuing a `CHANGE REPLICATION FILTER REPLICATE_REWRITE_DB` statement.

Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named *`channel_1`* use `--replicate-rewrite-db:`*`channel_1`*`:(`*`db_name1->db_name2`*`)`. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, “Replication Channel Based Filters” for more information.

Statements in which table names are qualified with database names when using this option do not work with table-level replication filtering options such as `--replicate-do-table`. Suppose we have a database named `a` on the master, one named `b` on the slave, each containing a table `t`, and have started the master with `--replicate-rewrite-db='a->b'`. At a later point in time, we execute `DELETE FROM a.t`. In this case, no relevant filtering rule works, for the reasons shown here:

1. `--replicate-do-table=a.t` does not work because the slave has table `t` in database `b`.
2. `--replicate-do-table=b.t` does not match the original statement and so is ignored.
3. `--replicate-do-table=*.t` is handled identically to `--replicate-do-table=a.t`, and thus does not work, either.

Similarly, the `--replication-rewrite-db` option does not work with cross-database updates.

If you use this option on the command line and the `>` character is special to your command interpreter, quote the option value. For example:

```
shell> mysqld --replicate-rewrite-db="olddb->newdb"
```

- `--replicate-same-server-id`

Command-Line Format	`--replicate-same-server-id`	
Permitted Values	Type	boolean
	Default	`FALSE`

To be used on slave servers. Usually you should use the default setting of 0, to prevent infinite loops caused by circular replication. If set to 1, the slave does not skip events having its own server ID. Normally, this is useful only in rare configurations. Cannot be set to 1 if `--log-slave-updates` is used. By default, the slave I/O thread does not write binary log events to the relay log if they have the slave's server ID (this optimization helps save disk usage). If you want to use `--replicate-same-server-id`, be sure to start the slave with this option before you make the slave read its own events that you want the slave SQL thread to execute.

- `--replicate-wild-do-table=`*`db_name.tbl_name`*

Command-Line Format	`--replicate-wild-do-table=name`	
Permitted Values	Type	string

Creates a replication filter by telling the slave thread to restrict replication to statements where any of the updated tables match the specified database and table name patterns. Patterns can contain

the % and _ wildcard characters, which have the same meaning as for the LIKE pattern-matching operator. To specify more than one table, use this option multiple times, once for each table. This works for cross-database updates. See Section 18.2.5, "How Servers Evaluate Replication Filtering Rules". You can also create such a filter by issuing a CHANGE REPLICATION FILTER REPLICATE_WILD_DO_TABLE statement.

Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named *channel_1* use --replicate-wild-do-table:*channel_1*:*db_name.tbl_name*. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, "Replication Channel Based Filters" for more information.

This option applies to tables, views, and triggers. It does not apply to stored procedures and functions, or events. To filter statements operating on the latter objects, use one or more of the --replicate-*-db options.

Example: --replicate-wild-do-table=foo%.bar% replicates only updates that use a table where the database name starts with foo and the table name starts with bar.

If the table name pattern is %, it matches any table name and the option also applies to database-level statements (CREATE DATABASE, DROP DATABASE, and ALTER DATABASE). For example, if you use --replicate-wild-do-table=foo%.%, database-level statements are replicated if the database name matches the pattern foo%.

To include literal wildcard characters in the database or table name patterns, escape them with a backslash. For example, to replicate all tables of a database that is named my_own%db, but not replicate tables from the my1ownAABCdb database, you should escape the _ and % characters like this: --replicate-wild-do-table=my_own\%db. If you use the option on the command line, you might need to double the backslashes or quote the option value, depending on your command interpreter. For example, with the bash shell, you would need to type --replicate-wild-do-table=my_own\\%db.

- --replicate-wild-ignore-table=*db_name.tbl_name*

Command-Line Format	--replicate-wild-ignore-table=name	
Permitted Values	Type	string

Creates a replication filter which keeps the slave thread from replicating a statement in which any table matches the given wildcard pattern. To specify more than one table to ignore, use this option multiple times, once for each table. This works for cross-database updates. See Section 18.2.5, "How Servers Evaluate Replication Filtering Rules". You can also create such a filter by issuing a CHANGE REPLICATION FILTER REPLICATE_WILD_IGNORE_TABLE statement.

Supports replication channel specific filters, enabling multi-source replication slaves to use specific filters for different sources. To configure a replication channel specific filter on a channel named *channel_1* use --replicate-wild-ignore:*channel_1*:*db_name.tbl_name*. In this case, the first colon is interpreted as a separator and subsequent colons are literal colons. See Section 18.2.5.4, "Replication Channel Based Filters" for more information.

Example: --replicate-wild-ignore-table=foo%.bar% does not replicate updates that use a table where the database name starts with foo and the table name starts with bar.

For information about how matching works, see the description of the --replicate-wild-do-table option. The rules for including literal wildcard characters in the option value are the same as for --replicate-wild-ignore-table as well.

- --report-host=*host_name*

Command-Line Format	--report-host=host_name

System Variable	Name	report_host
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The host name or IP address of the slave to be reported to the master during slave registration. This value appears in the output of `SHOW SLAVE HOSTS` on the master server. Leave the value unset if you do not want the slave to register itself with the master.

Note

It is not sufficient for the master to simply read the IP address of the slave from the TCP/IP socket after the slave connects. Due to NAT and other routing issues, that IP may not be valid for connecting to the slave from the master or other hosts.

- `--report-password=`*`password`*

Command-Line Format	--report-password=name	
System Variable	Name	report_password
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The account password of the slave to be reported to the master during slave registration. This value appears in the output of `SHOW SLAVE HOSTS` on the master server if the `--show-slave-auth-info` option is given.

Although the name of this option might imply otherwise, `--report-password` is not connected to the MySQL user privilege system and so is not necessarily (or even likely to be) the same as the password for the MySQL replication user account.

- `--report-port=`*`slave_port_num`*

Command-Line Format	--report-port=#	
System Variable	Name	report_port
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	[slave_port]
	Min Value	0
	Max Value	65535

The TCP/IP port number for connecting to the slave, to be reported to the master during slave registration. Set this only if the slave is listening on a nondefault port or if you have a special tunnel from the master or other clients to the slave. If you are not sure, do not use this option.

The default value for this option is the port number actually used by the slave (Bug #13333431). This is also the default value displayed by `SHOW SLAVE HOSTS`.

- `--report-user=user_name`

Command-Line Format	`--report-user=name`	
System Variable	Name	`report_user`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

The account user name of the slave to be reported to the master during slave registration. This value appears in the output of `SHOW SLAVE HOSTS` on the master server if the `--show-slave-auth-info` option is given.

Although the name of this option might imply otherwise, `--report-user` is not connected to the MySQL user privilege system and so is not necessarily (or even likely to be) the same as the name of the MySQL replication user account.

- `--show-slave-auth-info`

Command-Line Format	`--show-slave-auth-info`	
Permitted Values	Type	boolean
	Default	`FALSE`

Display slave user names and passwords in the output of `SHOW SLAVE HOSTS` on the master server for slaves started with the `--report-user` and `--report-password` options.

- `--slave-checkpoint-group=#`

Command-Line Format	`--slave-checkpoint-group=#`	
Permitted Values	Type	integer
	Default	`512`
	Min Value	`32`
	Max Value	`524280`
	Block Size	`8`

Sets the maximum number of transactions that can be processed by a multi-threaded slave before a checkpoint operation is called to update its status as shown by `SHOW SLAVE STATUS`. Setting this option has no effect on slaves for which multi-threading is not enabled.

This option works in combination with the `--slave-checkpoint-period` option in such a way that, when either limit is exceeded, the checkpoint is executed and the counters tracking both the number of transactions and the time elapsed since the last checkpoint are reset.

The minimum allowed value for this option is 32, unless the server was built using `-DWITH_DEBUG`, in which case the minimum value is 1. The effective value is always a multiple of 8; you can set it to a value that is not such a multiple, but the server rounds it down to the next lower multiple of 8 before storing the value. (*Exception*: No such rounding is performed by the debug server.) Regardless of how the server was built, the default value is 512, and the maximum allowed value is 524280.

- `--slave-checkpoint-period=#`

Command-Line Format	`--slave-checkpoint-period=#`	
Permitted Values	**Type**	integer
	Default	`300`
	Min Value	`1`
	Max Value	`4G`

Sets the maximum time (in milliseconds) that is allowed to pass before a checkpoint operation is called to update the status of a multi-threaded slave as shown by `SHOW SLAVE STATUS`. Setting this option has no effect on slaves for which multi-threading is not enabled.

This option works in combination with the `--slave-checkpoint-group` option in such a way that, when either limit is exceeded, the checkpoint is executed and the counters tracking both the number of transactions and the time elapsed since the last checkpoint are reset.

The minimum allowed value for this option is 1, unless the server was built using `-DWITH_DEBUG`, in which case the minimum value is 0. Regardless of how the server was built, the default value is 300, and the maximum possible value is 4294967296 (4GB).

- `--slave-parallel-workers`

Command-Line Format	`--slave-parallel-workers=#`	
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1024`

Sets the number of slave applier threads for executing replication transactions in parallel. Setting this variable to a number greater than 0 creates a multi-threaded slave with this number of applier threads. When set to 0 (the default) parallel execution is disabled and the slave uses a single applier thread.

A multi-threaded slave provides parallel execution by using a coordinator thread and the number of applier threads configured by this option. The way which transactions are distributed among applier threads is configured by `--slave-parallel-type`. For more information about multi-threaded slaves see `slave-parallel-workers`.

- `--slave-pending-jobs-size-max=#`

Command-Line Format	`--slave-pending-jobs-size-max=#`	
Permitted Values	**Type**	integer
	Default	`16M`

	Min Value	1024
	Max Value	18EB
	Block Size	1024

For multi-threaded slaves, this option sets the maximum amount of memory (in bytes) available to slave worker queues holding events not yet applied. Setting this option has no effect on slaves for which multi-threading is not enabled.

The minimum possible value for this option is 1024; the default is 16MB. The maximum possible value is 18446744073709551615 (16 exabytes). Values that are not exact multiples of 1024 are rounded down to the next-highest multiple of 1024 prior to being stored.

Important

The value for this option must not be less than the master's value for `max_allowed_packet`; otherwise a slave worker queue may become full while there remain events coming from the master to be processed.

- `--skip-slave-start`

Command-Line Format	`--skip-slave-start`	
Permitted Values	**Type**	boolean
	Default	`FALSE`

Tells the slave server not to start the slave threads when the server starts. To start the threads later, use a `START SLAVE` statement.

- `--slave_compressed_protocol={0|1}`

Command-Line Format	`--slave-compressed-protocol`	
System Variable	**Name**	`slave_compressed_protocol`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

If this option is set to 1, use compression for the slave/master protocol if both the slave and the master support it. The default is 0 (no compression).

- `--slave-load-tmpdir=`*`dir_name`*

Command-Line Format	`--slave-load-tmpdir=dir_name`	
System Variable	**Name**	`slave_load_tmpdir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

	Default	`/tmp`

The name of the directory where the slave creates temporary files. This option is by default equal to the value of the `tmpdir` system variable. When the slave SQL thread replicates a `LOAD DATA INFILE` statement, it extracts the file to be loaded from the relay log into temporary files, and then loads these into the table. If the file loaded on the master is huge, the temporary files on the slave are huge, too. Therefore, it might be advisable to use this option to tell the slave to put temporary files in a directory located in some file system that has a lot of available space. In that case, the relay logs are huge as well, so you might also want to use the `--relay-log` option to place the relay logs in that file system.

The directory specified by this option should be located in a disk-based file system (not a memory-based file system) because the temporary files used to replicate `LOAD DATA INFILE` must survive machine restarts. The directory also should not be one that is cleared by the operating system during the system startup process.

- `slave-max-allowed-packet=`*`bytes`*

Command-Line Format	`--slave-max-allowed-packet=#`	
Permitted Values	**Type**	integer
	Default	`1073741824`
	Min Value	`1024`
	Max Value	`1073741824`

This option sets the maximum packet size in bytes for the slave SQL and I/O threads, so that large updates using row-based replication do not cause replication to fail because an update exceeded `max_allowed_packet`. (Bug #12400221, Bug #60926)

The corresponding server variable `slave_max_allowed_packet` always has a value that is a positive integer multiple of 1024; if you set it to some value that is not such a multiple, the value is automatically rounded down to the next highest multiple of 1024. (For example, if you start the server with `--slave-max-allowed-packet=10000`, the value used is 9216; setting 0 as the value causes 1024 to be used.) A truncation warning is issued in such cases.

The maximum (and default) value is 1073741824 (1 GB); the minimum is 1024.

- `--slave-net-timeout=`*`seconds`*

Command-Line Format	`--slave-net-timeout=#`	
System Variable	**Name**	`slave_net_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`3600`
	Min Value	`1`
Permitted Values	**Type**	integer
	Default	`60`

	Min Value	1

The number of seconds to wait for more data from the master before the slave considers the connection broken, aborts the read, and tries to reconnect. The first retry occurs immediately after the timeout. The interval between retries is controlled by the MASTER_CONNECT_RETRY option for the CHANGE MASTER TO statement, and the number of reconnection attempts is limited by the --master-retry-count option. The default value is 60 seconds (one minute).

- --slave-parallel-type=*type*

Command-Line Format	--slave-parallel-type=type	
Permitted Values	**Type**	enumeration
	Default	DATABASE
	Valid Values	DATABASE
		LOGICAL_CLOCK

When using a multi-threaded slave (slave_parallel_workers is greater than 0), this option specifies the policy used to decide which transactions are allowed to execute in parallel on the slave. The possible values are:

- DATABASE: Transactions that update different databases are applied in parallel. This value is only appropriate if data is partitioned into multiple databases which are being updated independently and concurrently on the master. Only recommended if there are no cross-database constraints, as such constraints may be violated on the slave.
- LOGICAL_CLOCK: Transactions that are part of the same binary log group commit on a master are applied in parallel on a slave. There are no cross-database constraints, and data does not need to be partitioned into multiple databases.

Regardless of the value of this variable, there is no special configuration required on the master. When slave_preserve_commit_order=1, you can only use LOGICAL_CLOCK. If your replication topology uses multiple levels of slaves, LOGICAL_CLOCK may achieve less parallelization for each level the slave is away from the master.

- slave-rows-search-algorithms=*list*

Command-Line Format	--slave-rows-search-algorithms=list	
Permitted Values	**Type**	set
	Default	TABLE_SCAN,INDEX_SCAN
	Valid Values	TABLE_SCAN,INDEX_SCAN
		INDEX_SCAN,HASH_SCAN
		TABLE_SCAN,HASH_SCAN
		TABLE_SCAN,INDEX_SCAN,HASH_SCAN (equivalent to INDEX_SCAN,HASH_SCAN)
Permitted Values (>= 8.0.2)	**Type**	set
	Default	INDEX_SCAN,HASH_SCAN
	Valid Values	TABLE_SCAN,INDEX_SCAN
		INDEX_SCAN,HASH_SCAN
		TABLE_SCAN,HASH_SCAN
		TABLE_SCAN,INDEX_SCAN,HASH_SCAN (equivalent to INDEX_SCAN,HASH_SCAN)

When preparing batches of rows for row-based logging and replication, this option controls how the rows are searched for matches—that is, whether or not hashing is used for searches using a primary or unique key, some other key, or no key at all. This option takes a comma-separated list of any 2 (or possibly 3) values from the list `INDEX_SCAN`, `TABLE_SCAN`, `HASH_SCAN`. The list need not be quoted, but must contain no spaces, whether or not quotes are used. Possible combinations (lists) and their effects are shown in the following table:

Index used / option value	`INDEX_SCAN,HASH_SCAN` or `INDEX_SCAN,TABLE_SCAN,HASH_SCAN`	`INDEX_SCAN,TABLE_SCAN`	`TABLE_SCAN,HASH_SCAN`
Primary key or unique key	Index scan	Index scan	Hash scan over index
(Other) Key	Hash scan over index	Index scan	Hash scan over index
No index	Hash scan	Table scan	Hash scan

The order in which the algorithms are specified in the list does not make any difference in the order in which they are displayed by a `SELECT` or `SHOW VARIABLES` statement (which is the same as that used in the table just shown previously).The default value is `TABLE_SCAN,INDEX_SCAN`, which means that all searches that can use indexes do use them, and searches without any indexes use table scans.

Specifying `INDEX_SCAN,TABLE_SCAN,HASH_SCAN` has the same effect as specifying `INDEX_SCAN,HASH_SCAN`. To use hashing for any searches that does not use a primary or unique key, set this option to `INDEX_SCAN,HASH_SCAN`. To force hashing for *all* searches, set it to `TABLE_SCAN,HASH_SCAN`.

Note

There is only a performance advantage for `INDEX_SCAN` and `HASH_SCAN` if the row events are big enough. The size of row events is configured using `--binlog-row-event-max-size`. For example, suppose a `DELETE` statement which deletes 25,000 rows generates large `Delete_row_event` events. In this case if `slave_rows_search_algorithms` is set to `INDEX_SCAN` or `HASH_SCAN` there is a performance improvement. However, if there are 25,000 `DELETE` statements and each is represented by a separate event then setting `slave_rows_search_algorithms` to `INDEX_SCAN` or `HASH_SCAN` provides no performance improvement while executing these separate events.

- `--slave-skip-errors=[err_code1,err_code2,...|all|ddl_exist_errors]`

Command-Line Format	`--slave-skip-errors=name`	
System Variable	**Name**	`slave_skip_errors`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`OFF`
	Valid Values	`OFF`
		`[list of error codes]`
		`all`
		`ddl_exist_errors`

Permitted Values	**Type**	string
	Default	`OFF`
	Valid Values	`OFF`
		`[list of error codes]`
		`all`
		`ddl_exist_errors`
Permitted Values	**Type**	string
	Default	`OFF`
	Valid Values	`OFF`
		`[list of error codes]`
		`all`
		`ddl_exist_errors`

Normally, replication stops when an error occurs on the slave, which gives you the opportunity to resolve the inconsistency in the data manually. This option causes the slave SQL thread to continue replication when a statement returns any of the errors listed in the option value.

Do not use this option unless you fully understand why you are getting errors. If there are no bugs in your replication setup and client programs, and no bugs in MySQL itself, an error that stops replication should never occur. Indiscriminate use of this option results in slaves becoming hopelessly out of synchrony with the master, with you having no idea why this has occurred.

For error codes, you should use the numbers provided by the error message in your slave error log and in the output of `SHOW SLAVE STATUS`. Appendix B, *Errors, Error Codes, and Common Problems*, lists server error codes.

You can also (but should not) use the very nonrecommended value of `all` to cause the slave to ignore all error messages and keeps going regardless of what happens. Needless to say, if you use `all`, there are no guarantees regarding the integrity of your data. Please do not complain (or file bug reports) in this case if the slave's data is not anywhere close to what it is on the master. *You have been warned*.

MySQL 8.0 supports an additional shorthand value `ddl_exist_errors`, which is equivalent to the error code list `1007,1008,1050,1051,1054,1060,1061,1068,1094,1146`.

Examples:

```
--slave-skip-errors=1062,1053
--slave-skip-errors=all
--slave-skip-errors=ddl_exist_errors
```

- `--slave-sql-verify-checksum={0|1}`

Command-Line Format	`--slave-sql-verify-checksum=value`	
Permitted Values	**Type**	boolean
	Default	`0`
	Valid Values	`0`
		`1`

When this option is enabled, the slave examines checksums read from the relay log, in the event of a mismatch, the slave stops with an error. Disabled by default.

The following options are used internally by the MySQL test suite for replication testing and debugging. They are not intended for use in a production setting.

- `--abort-slave-event-count`

Command-Line Format	`--abort-slave-event-count=#`	
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`

When this option is set to some positive integer *`value`* other than 0 (the default) it affects replication behavior as follows: After the slave SQL thread has started, *`value`* log events are permitted to be executed; after that, the slave SQL thread does not receive any more events, just as if the network connection from the master were cut. The slave thread continues to run, and the output from `SHOW SLAVE STATUS` displays `Yes` in both the `Slave_IO_Running` and the `Slave_SQL_Running` columns, but no further events are read from the relay log.

- `--disconnect-slave-event-count`

Command-Line Format	`--disconnect-slave-event-count=#`	
Permitted Values	Type	integer
	Default	`0`

Options for Logging Slave Status to Tables

MySQL 8.0 supports logging of replication slave status information to tables rather than files. Writing of the master info log and the relay log info log can be configured separately using the two server options listed here:

- `--master-info-repository={FILE|TABLE}`

Command-Line Format	`--master-info-repository=FILE|TABLE`	
Permitted Values (<= 8.0.1)	Type	string
	Default	`FILE`
	Valid Values	`FILE`
		`TABLE`
Permitted Values (>= 8.0.2)	Type	string
	Default	`TABLE`
	Valid Values	`FILE`
		`TABLE`

This option causes the server to write its master info log to a file or a table. The name of the file defaults to `master.info`; you can change the name of the file using the `--master-info-file` server option.

The default value for this option is `FILE`. If you use `TABLE`, the log is written to the `slave_master_info` table in the `mysql` database.

- `--relay-log-info-repository={FILE|TABLE}`

Command-Line Format	`--relay-log-info-repository=FILE|TABLE`	
Permitted Values	Type	string
	Default	`FILE`

	Valid Values	`FILE`
		`TABLE`

This option causes the server to log its relay log info to a file or a table. The name of the file defaults to `relay-log.info`; you can change the name of the file using the `--relay-log-info-file` server option.

The default value for this option is `FILE`. If you use `TABLE`, the log is written to the `slave_relay_log_info` table in the `mysql` database.

For replication to be crash-safe, this option must be set to `TABLE`; in additon, the `--relay-log-recovery` option must be enabled. See Crash-safe replication, for more information.

The info log tables and their contents are considered local to a given MySQL Server. They are not replicated, and changes to them are not written to the binary log.

For more information, see Section 18.2.4, “Replication Relay and Status Logs”.

System Variables Used on Replication Slaves

The following list describes system variables for controlling replication slave servers. They can be set at server startup and some of them can be changed at runtime using `SET`. Server options used with replication slaves are listed earlier in this section.

- `slave_allow_batching`

Command-Line Format	`--slave-allow-batching`	
System Variable	**Name**	`slave_allow_batching`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`off`

Whether or not batched updates are enabled on replication slaves.

Setting this variable has an effect only when using replication with the `NDB` storage engine; in MySQL Server 8.0, it is present but does nothing. For more information, see Starting NDB Cluster Replication (Single Replication Channel).

- `init_slave`

Command-Line Format	`--init-slave=name`	
System Variable	**Name**	`init_slave`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

This variable is similar to `init_connect`, but is a string to be executed by a slave server each time the SQL thread starts. The format of the string is the same as for the `init_connect` variable. The setting of this variable takes effect for subsequent `START SLAVE` statements.

Note

The SQL thread sends an acknowledgment to the client before it executes `init_slave`. Therefore, it is not guaranteed that `init_slave` has been executed when `START SLAVE` returns. See Section 13.4.2.6, "START SLAVE Syntax", for more information.

- `log_slow_slave_statements`

System Variable	Name	log_slow_slave_statements
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

When the slow query log is enabled, this variable enables logging for queries that have taken more than `long_query_time` seconds to execute on the slave. Setting this variable has no immediate effect. The state of the variable applies on all subsequent `START SLAVE` statements.

- `master_info_repository`

Command-Line Format	--master-info-repository=FILE|TABLE	
System Variable	Name	master_info_repository
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (<= 8.0.1)	Type	string
	Default	FILE
	Valid Values	FILE
		TABLE
Permitted Values (>= 8.0.2)	Type	string
	Default	TABLE
	Valid Values	FILE
		TABLE

The setting of this variable determines whether the slave logs master status and connection information to a `FILE` (`master.info`), or to a `TABLE` (`mysql.slave_master_info`). You can only change the value of this variable when no replication threads are executing.

The setting of this variable also has a direct influence on the effect had by the setting of the `sync_master_info` system variable; see that variable's description for further information.

This variable must be set to `TABLE` before configuring multiple replication channels. If you are using multiple replication channels then you cannot set this variable back to `FILE`.

- `relay_log`

Command-Line Format	--relay-log=file_name	
System Variable	Name	relay_log

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The base name of the relay log file, with no paths and no file extension. By default `relay-log`. The file name of individual files for the default replication channel is `relay-log.XXXXXX`, and for additional replication channels is `relay-log-channel.XXXXXX`.

- `relay_log_basename`

System Variable	**Name**	`relay_log_basename`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name
	Default	`datadir + '/' + hostname + '-relay-bin'`

Holds the name and complete path to the relay log file.

- `relay_log_index`

Command-Line Format	`--relay-log-index`	
System Variable	**Name**	`relay_log_index`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name
	Default	`*host_name*-relay-bin.index`

The name of the relay log index file for the default replication channel. The default name is *`host_name`*`-relay-bin.index` in the data directory, where *`host_name`* is the name of the slave server.

- `relay_log_info_file`

Command-Line Format	`--relay-log-info-file=file_name`	
System Variable	**Name**	`relay_log_info_file`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name
	Default	`relay-log.info`

The name of the file in which the slave records information about the relay logs, when `relay_log_info_repository=FILE`. If `relay_log_info_repository=TABLE`, it is the file name that would be used in case the repository was changed to `FILE`). The default name is `relay-log.info` in the data directory.

- `relay_log_info_repository`

System Variable	Name	relay_log_info_repository
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string
	Default	FILE
	Valid Values	FILE
		TABLE

This variable determines whether the slave's position in the relay logs is written to a `FILE` (`relay-log.info`) or to a `TABLE` (`mysql.slave_relay_log_info`). You can only change the value of this variable when no replication threads are executing.

The setting of this variable also has a direct influence on the effect had by the setting of the `sync_relay_log_info` system variable; see that variable's description for further information.

This variable must be set to `TABLE` before configuring multiple replication channels. If you are using multiple replication channels then you cannot set this variable back to `FILE`.

- `relay_log_recovery`

Command-Line Format	--relay-log-recovery	
System Variable	Name	relay_log_recovery
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	FALSE

Enables automatic relay log recovery immediately following server startup. The recovery process creates a new relay log file, initializes the SQL thread position to this new relay log, and initializes the I/O thread to the SQL thread position. Reading of the relay log from the master then continues. In MySQL 5.7, this global variable is read-only; its value can be changed by starting the slave with the `--relay-log-recovery` option, which should be used following a crash on the replication slave to ensure that no possibly corrupted relay logs are processed, and must be used in order to guarantee a crash-safe slave. The default value is 0 (disabled).

This variable also interacts with `relay-log-purge`, which controls purging of logs when they are no longer needed. Enabling the `--relay-log-recovery` option when `relay-log-purge` is disabled risks reading the relay log from files that were not purged, leading to data inconsistency, and is therefore not crash-safe.

When `relay_log_recovery` is enabled and the slave has stopped due to errors encountered while running in multi-threaded mode, you can use `START SLAVE UNTIL SQL_AFTER_MTS_GAPS` to ensure that all gaps are processed before switching back to single-threaded mode or executing a `CHANGE MASTER TO` statement.

- `rpl_stop_slave_timeout`

Command-Line Format	--rpl-stop-slave-timeout=seconds

System Variable	Name	rpl_stop_slave_timeout
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	31536000
	Min Value	2
	Max Value	31536000

You can control the length of time (in seconds) that `STOP SLAVE` waits before timing out by setting this variable. This can be used to avoid deadlocks between `STOP SLAVE` and other slave SQL statements using different client connections to the slave. The maximum and default value of `rpl_stop_slave_timeout` is 31536000 seconds (1 year). The minimum is 2 seconds. Changes to this variable take effect for subsequent `STOP SLAVE` statements. This variable affects only the client that issues a `STOP SLAVE` statement. When the timeout is reached, the issuing client stops waiting for the slave threads to stop, but the slave threads continue to try to stop.

- `slave_checkpoint_group`

Command-Line Format	--slave-checkpoint-group=#	
System Variable	Name	slave_checkpoint_group=#
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	512
	Min Value	32
	Max Value	524280
	Block Size	8

Sets the maximum number of transactions that can be processed by a multi-threaded slave before a checkpoint operation is called to update its status as shown by `SHOW SLAVE STATUS`. Setting this variable has no effect on slaves for which multi-threading is not enabled. Setting this variable has no immediate effect. The state of the variable applies on all subsequent `START SLAVE` commands.

This variable works in combination with the `slave_checkpoint_period` system variable in such a way that, when either limit is exceeded, the checkpoint is executed and the counters tracking both the number of transactions and the time elapsed since the last checkpoint are reset.

The minimum allowed value for this variable is 32, unless the server was built using `-DWITH_DEBUG`, in which case the minimum value is 1. The effective value is always a multiple of 8; you can set it to a value that is not such a multiple, but the server rounds it down to the next lower multiple of 8 before storing the value. (*Exception*: No such rounding is performed by the debug server.) Regardless of how the server was built, the default value is 512, and the maximum allowed value is 524280.

- `slave_checkpoint_period`

Command-Line Format	--slave-checkpoint-period=#	
System Variable	Name	slave_checkpoint_period=#
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	300
	Min Value	1
	Max Value	4G

Sets the maximum time (in milliseconds) that is allowed to pass before a checkpoint operation is called to update the status of a multi-threaded slave as shown by SHOW SLAVE STATUS. Setting this variable has no effect on slaves for which multi-threading is not enabled. Setting this variable takes effect for all replication channels immediately, including running channels.

This variable works in combination with the slave_checkpoint_group system variable in such a way that, when either limit is exceeded, the checkpoint is executed and the counters tracking both the number of transactions and the time elapsed since the last checkpoint are reset.

The minimum allowed value for this variable is 1, unless the server was built using -DWITH_DEBUG, in which case the minimum value is 0. Regardless of how the server was built, the default value is 300, and the maximum possible value is 4294967296 (4GB).

- slave_compressed_protocol

Command-Line Format	--slave-compressed-protocol	
System Variable	Name	slave_compressed_protocol
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	OFF

Whether to use compression of the slave/master protocol if both the slave and the master support it. Changes to this variable take effect on subsequent connection attempts; this includes after issuing a START SLAVE statement, as well as reconnections made by a running I/O thread (for example after issuing a CHANGE MASTER TO MASTER_RETRY_COUNT statement).

- slave_exec_mode

Command-Line Format	--slave-exec-mode=mode	
System Variable	Name	slave_exec_mode
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration

	Default	`STRICT` (ALL)
	Default	`IDEMPOTENT` (NDB)
	Valid Values	`IDEMPOTENT`
		`STRICT`

Controls how a slave thread resolves conflicts and errors during replication. `IDEMPOTENT` mode causes suppression of duplicate-key and no-key-found errors. This mode should be employed in multi-master replication, circular replication, and some other special replication scenarios. `STRICT` mode is the default, and is suitable for most other cases. Setting this variable takes effect for all replication channels immediately, including running channels.

- `slave_load_tmpdir`

Command-Line Format	`--slave-load-tmpdir=dir_name`	
System Variable	**Name**	`slave_load_tmpdir`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name
	Default	`/tmp`

The name of the directory where the slave creates temporary files for replicating `LOAD DATA INFILE` statements. Setting this variable takes effect for all replication channels immediately, including running channels.

- `slave_max_allowed_packet`

System Variable	**Name**	`slave_max_allowed_packet`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1073741824`
	Min Value	`1024`
	Max Value	`1073741824`

This variable sets the maximum packet size for the slave SQL and I/O threads, so that large updates using row-based replication do not cause replication to fail because an update exceeded `max_allowed_packet`. Setting this variable takes effect for all replication channels immediately, including running channels.

This global variable always has a value that is a positive integer multiple of 1024; if you set it to some value that is not, the value is rounded down to the next highest multiple of 1024 for it is stored or used; setting `slave_max_allowed_packet` to 0 causes 1024 to be used. (A truncation warning is issued in all such cases.) The default and maximum value is 1073741824 (1 GB); the minimum is 1024.

`slave_max_allowed_packet` can also be set at startup, using the `--slave-max-allowed-packet` option.

- `slave_net_timeout`

Command-Line Format		`--slave-net-timeout=#`
System Variable	Name	`slave_net_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`3600`
	Min Value	`1`
Permitted Values	Type	integer
	Default	`60`
	Min Value	`1`

The number of seconds to wait for more data from a master/slave connection before aborting the read. Setting this variable has no immediate effect. The state of the variable applies on all subsequent `START SLAVE` commands.

- `slave_parallel_type=type`

Command-Line Format		`--slave-parallel-type=type`
Permitted Values	Type	enumeration
	Default	`DATABASE`
	Valid Values	`DATABASE`
		`LOGICAL_CLOCK`

When using a multi-threaded slave (`slave_parallel_workers` is greater than 0), this variable specifies the policy used to decide which transactions are allowed to execute in parallel on the slave. See `--slave-parallel-type` for more information.

- `slave_parallel_workers`

Command-Line Format		`--slave-parallel-workers=#`
System Variable	Name	`slave_parallel_workers`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1024`

Sets the number of slave applier threads for executing replication transactions in parallel. Setting this variable to a number greater than 0 creates a multi-threaded slave with this number of applier threads. When set to 0 (the default) parallel execution is disabled and the slave uses a single applier

thread. Setting `slave_parallel_workers` has no immediate effect. The state of the variable applies on all subsequent `START SLAVE` statements. Retrying of transactions is supported when multi-threading is enabled on a slave.

A multi-threaded slave provides parallel execution by using a coordinator thread and the number of applier threads configured by this variable. The way which transactions are distributed among applier threads is configured by `slave_parallel_type`. The transactions that the slave applies in parallel may commit out of order, unless `slave_preserve_commit_order=1`. Therefore, checking for the most recently executed transaction does not guarantee that all previous transactions from the master have been executed on the slave. This has implications for logging and recovery when using a multi-threaded slave. For example, on a multi-threaded slave the `START SLAVE UNTIL` statement only supports using `SQL_AFTER_MTS_GAPS`.

- `slave_pending_jobs_size_max`

System Variable	**Name**	`slave_pending_jobs_size_max`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`16M`
	Min Value	`1024`
	Max Value	`18EB`
	Block Size	`1024`

For multi-threaded slaves, this variable sets the maximum amount of memory (in bytes) available to slave worker queues holding events not yet applied. Setting this variable has no effect on slaves for which multi-threading is not enabled. Setting this variable has no immediate effect. The state of the variable applies on all subsequent `START SLAVE` commands.

The minimum possible value for this variable is 1024; the default is 16MB. The maximum possible value is 18446744073709551615 (16 exabytes). Values that are not exact multiples of 1024 are rounded down to the next-highest multiple of 1024 prior to being stored.

Important

The value of this variable must not be less than the master's value for `max_allowed_packet`; otherwise a slave worker queue may become full while there remain events coming from the master to be processed.

- `slave_preserve_commit_order`

Command-Line Format	`--slave-preserve-commit-order=value`	
System Variable	**Name**	`slave_preserve_commit_order`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`0`

	Valid Values	`0`
		`1`

For multi-threaded slaves, enabling this variable ensures that transactions are externalized on the slave in the same order as they appear in the slave's relay log. Setting this variable has no effect on slaves for which multi-threading is not enabled. All replication threads (for all replication channels if you are using multiple replication channels) must be stopped before changing this variable. `--log-bin` and `--log-slave-updates` must be enabled on the slave. In addition `--slave-parallel-type` must be set to `LOGICAL_CLOCK`.

Once a multi-threaded slave has been started, transactions can begin to execute in parallel. With `slave_preserve_commit_order` enabled, the executing thread waits until all previous transactions are committed before committing. While the slave thread is waiting for other workers to commit their transactions it reports its status as `Waiting for preceding transaction to commit`. Enabling this mode on a multi-threaded slave ensures that it never enters a state that the master was not in. This makes it well suited to using replication for read scale-out. See Section 18.3.5, “Using Replication for Scale-Out”.

When using a multi-threaded slave, if `slave_preserve_commit_order` is not enabled, there is a chance of gaps in the sequence of transactions that have been executed from the slave's relay log. When this option is enabled, there is not this chance of gaps, but `Exec_master_log_pos` may be behind the position up to which has been executed. See Section 18.4.1.34, “Replication and Transaction Inconsistencies” for more information.

- `slave_rows_search_algorithms`

System Variable	**Name**	`slave_rows_search_algorithms=list`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	set
	Default	`TABLE_SCAN,INDEX_SCAN`
	Valid Values	`TABLE_SCAN,INDEX_SCAN`
		`INDEX_SCAN,HASH_SCAN`
		`TABLE_SCAN,HASH_SCAN`
		`TABLE_SCAN,INDEX_SCAN,HASH_SCAN` (equivalent to INDEX_SCAN,HASH_SCAN)
Permitted Values (>= 8.0.2)	**Type**	set
	Default	`INDEX_SCAN,HASH_SCAN`
	Valid Values	`TABLE_SCAN,INDEX_SCAN`
		`INDEX_SCAN,HASH_SCAN`
		`TABLE_SCAN,HASH_SCAN`
		`TABLE_SCAN,INDEX_SCAN,HASH_SCAN` (equivalent to INDEX_SCAN,HASH_SCAN)

When preparing batches of rows for row-based logging and replication, this variable controls how the rows are searched for matches—that is, whether or not hashing is used for searches using a primary or unique key, using some other key, or using no key at all. Setting this variable takes effect for all replication channels immediately, including running channels.

This variable takes a comma-separated list of at least 2 values from the list `INDEX_SCAN`, `TABLE_SCAN`, `HASH_SCAN`. The value expected as a string, so the value must be quoted. In addition,

the value must not contain any spaces. Possible combinations (lists) and their effects are shown in the following table:

Index used / option value	`INDEX_SCAN,HASH_SCAN` or `INDEX_SCAN,TABLE_SCAN,HASH_SCAN`	`INDEX_SCAN,TABLE_SCAN`	`TABLE_SCAN,HASH_SCAN`
Primary key or unique key	Index scan	index scan	Index hash
(Other) Key	Index hash	Index scan	Index hash
No index	Table hash	Table scan	Table hash

The order in which the algorithms are specified in the list does not make any difference in the order in which they are displayed by a `SELECT` or `SHOW VARIABLES` statement, as shown here:

```
mysql> SET GLOBAL slave_rows_search_algorithms = "INDEX_SCAN,TABLE_SCAN";
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW VARIABLES LIKE '%algorithms%';
+------------------------------+----------------------+
| Variable_name                | Value                |
+------------------------------+----------------------+
| slave_rows_search_algorithms | TABLE_SCAN,INDEX_SCAN |
+------------------------------+----------------------+
1 row in set (0.00 sec)

mysql> SET GLOBAL slave_rows_search_algorithms = "TABLE_SCAN,INDEX_SCAN";
Query OK, 0 rows affected (0.00 sec)

mysql> SHOW VARIABLES LIKE '%algorithms%';
+------------------------------+----------------------+
| Variable_name                | Value                |
+------------------------------+----------------------+
| slave_rows_search_algorithms | TABLE_SCAN,INDEX_SCAN |
+------------------------------+----------------------+
1 row in set (0.00 sec)
```

The default value is `TABLE_SCAN,INDEX_SCAN`, which means that all searches that can use indexes do use them, and searches without any indexes use table scans.

Specifying `INDEX_SCAN,TABLE_SCAN,HASH_SCAN` has the same effect as specifying `INDEX_SCAN,HASH_SCAN`. To use hashing for any searches that does not use a primary or unique key, set this variable to `INDEX_SCAN,HASH_SCAN`. To force hashing for *all* searches, set it to `TABLE_SCAN,HASH_SCAN`.

- `slave_skip_errors`

Command-Line Format	`--slave-skip-errors=name`	
System Variable	**Name**	`slave_skip_errors`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	string
	Default	`OFF`
	Valid Values	`OFF`
		`[list of error codes]`
		`all`

		ddl_exist_errors
Permitted Values	**Type**	string
	Default	OFF
	Valid Values	OFF
		[list of error codes]
		all
		ddl_exist_errors
Permitted Values	**Type**	string
	Default	OFF
	Valid Values	OFF
		[list of error codes]
		all
		ddl_exist_errors

Normally, replication stops when an error occurs on the slave, which gives you the opportunity to resolve the inconsistency in the data manually. This variable causes the slave SQL thread to continue replication when a statement returns any of the errors listed in the variable value. The setting of this variable takes effect immediately, even for running replication threads.

- slave_sql_verify_checksum

System Variable	**Name**	slave_sql_verify_checksum
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	1
	Valid Values	0
		1

Cause the slave SQL thread to verify data using the checksums read from the relay log. In the event of a mismatch, the slave stops with an error. Setting this variable takes effect for all replication channels immediately, including running channels.

Note

The slave I/O thread always reads checksums if possible when accepting events from over the network.

- slave_transaction_retries

Command-Line Format		--slave-transaction-retries=#
System Variable	**Name**	slave_transaction_retries
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	10

	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`10`
	Min Value	`0`
	Max Value	`18446744073709551615`

If a replication slave SQL thread fails to execute a transaction because of an `InnoDB` deadlock or because the transaction's execution time exceeded `InnoDB`'s `innodb_lock_wait_timeout`, it automatically retries `slave_transaction_retries` times before stopping with an error. The default value is 10. Setting this variable takes effect for all replication channels immediately, including running channels.

Retrying of transactions is supported when multi-threading is enabled on a slave. In previous versions, `slave_transaction_retries` was treated as equal to 0 when using multi-threaded slaves.

- `slave_type_conversions`

Command-Line Format	`--slave-type-conversions=set`	
System Variable	**Name**	`slave_type_conversions`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	set
	Default	
	Valid Values	`ALL_LOSSY`
		`ALL_NON_LOSSY`
		`ALL_SIGNED`
		`ALL_UNSIGNED`

Controls the type conversion mode in effect on the slave when using row-based replication. Its value is a comma-delimited set of zero or more elements from the list: `ALL_LOSSY`, `ALL_NON_LOSSY`, `ALL_SIGNED`, `ALL_UNSIGNED`. Set this variable to an empty string to disallow type conversions between the master and the slave. Setting this variable takes effect for all replication channels immediately, including running channels.

For additional information on type conversion modes applicable to attribute promotion and demotion in row-based replication, see Row-based replication: attribute promotion and demotion.

- `sql_slave_skip_counter`

System Variable	**Name**	`sql_slave_skip_counter`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	Type	integer

The number of events from the master that a slave server should skip. Setting the option has no immediate effect. The variable applies to the next `START SLAVE` statement; the next `START SLAVE` statement also changes the value back to 0. When this variable is set to a non-zero value and there are multiple replication channels configured, the `START SLAVE` statement can only be used with the `FOR CHANNEL channel` clause.

This option is incompatible with GTID-based replication, and must not be set to a nonzero value when `--gtid-mode=ON`. If you need to skip transactions when employing GTIDs, use `gtid_executed` from the master instead. See Injecting empty transactions, for information about how to do this.

Important

If skipping the number of events specified by setting this variable would cause the slave to begin in the middle of an event group, the slave continues to skip until it finds the beginning of the next event group and begins from that point. For more information, see Section 13.4.2.5, "SET GLOBAL sql_slave_skip_counter Syntax".

- `sync_master_info`

Command-Line Format		`--sync-master-info=#`
System Variable	Name	`sync_master_info`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	`10000`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	Type	integer
	Default	`10000`
	Min Value	`0`
	Max Value	`18446744073709551615`

The effects of this variable on a replication slave depend on whether the slave's `master_info_repository` is set to `FILE` or `TABLE`, as explained in the following paragraphs.

master_info_repository = FILE. If the value of `sync_master_info` is greater than 0, the slave synchronizes its `master.info` file to disk (using `fdatasync()`) after every `sync_master_info` events. If it is 0, the MySQL server performs no synchronization of the `master.info` file to disk; instead, the server relies on the operating system to flush its contents periodically as with any other file.

master_info_repository = TABLE. If the value of `sync_master_info` is greater than 0, the slave updates its master info repository table after every `sync_master_info` events. If it is 0, the table is never updated.

The default value for `sync_master_info` is 10000. Setting this variable takes effect for all replication channels immediately, including running channels.

- `sync_relay_log`

Command-Line Format	`--sync-relay-log=#`	
System Variable	**Name**	`sync_relay_log`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`10000`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`10000`
	Min Value	`0`
	Max Value	`18446744073709551615`

If the value of this variable is greater than 0, the MySQL server synchronizes its relay log to disk (using `fdatasync()`) after every `sync_relay_log` events are written to the relay log. Setting this variable takes effect for all replication channels immediately, including running channels.

Setting `sync_relay_log` to 0 causes no synchronization to be done to disk; in this case, the server relies on the operating system to flush the relay log's contents from time to time as for any other file.

A value of 1 is the safest choice because in the event of a crash you lose at most one event from the relay log. However, it is also the slowest choice (unless the disk has a battery-backed cache, which makes synchronization very fast).

- `sync_relay_log_info`

Command-Line Format	`--sync-relay-log-info=#`	
System Variable	**Name**	`sync_relay_log_info`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`10000`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer

	Default	10000
	Min Value	0
	Max Value	18446744073709551615

The effects of this variable on the slave depend on the server's relay_log_info_repository setting (FILE or TABLE), and if this is TABLE, additionally on whether the storage engine used by the relay log info table is transactional (such as InnoDB) or not (MyISAM). The effects of these factors on the behavior of the server for sync_relay_log_info values of zero and greater than zero are shown in the following table:

sync_relay_log_info	**relay_log_info_repository**		
	FILE	**TABLE**	
		Transactional	**Nontransactional**
N > 0	The slave synchronizes its relay-log.info file to disk (using fdatasync()) after every *N* transactions.	The table is updated after each transaction. (*N* is effectively ignored.)	The table is updated after every *N* events.
0	The MySQL server performs no synchronization of the relay-log.info file to disk; instead, the server relies on the operating system to flush its contents periodically as with any other file.		The table is never updated.

The default value for sync_relay_log_info is 10000. Setting this variable takes effect for all replication channels immediately, including running channels.

18.1.6.4 Binary Logging Options and Variables

Startup Options Used with Binary Logging

System Variables Used with Binary Logging

You can use the mysqld options and system variables that are described in this section to affect the operation of the binary log as well as to control which statements are written to the binary log. For additional information about the binary log, see Section 5.4.4, “The Binary Log”. For additional information about using MySQL server options and system variables, see Section 5.1.4, “Server Command Options”, and Section 5.1.5, “Server System Variables”.

Startup Options Used with Binary Logging

The following list describes startup options for enabling and configuring the binary log. System variables used with binary logging are discussed later in this section.

- --binlog-row-event-max-size=*N*

Command-Line Format	--binlog-row-event-max-size=#	
Permitted Values (32-bit platforms)	**Type**	integer
	Default	8192
	Min Value	256

	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`8192`
	Min Value	`256`
	Max Value	`18446744073709551615`

Specify the maximum size of a row-based binary log event, in bytes. Rows are grouped into events smaller than this size if possible. The value should be a multiple of 256. The default is 8192. See Section 18.2.1, “Replication Formats”.

- `--log-bin[=base_name]`

Command-Line Format	`--log-bin`	
System Variable	**Name**	`log_bin`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

Enable binary logging. The server logs all statements that change data to the binary log, which is used for backup and replication. See Section 5.4.4, “The Binary Log”.

The option value, if given, is the base name for the log sequence. The server creates binary log files in sequence by adding a numeric suffix to the base name. It is recommended that you specify a base name (see Section B.5.7, “Known Issues in MySQL”, for the reason). Otherwise, MySQL uses *`host_name`*`-bin` as the base name.

When the server reads an entry from the index file, it checks whether the entry contains a relative path, and if it does, the relative part of the path in replaced with the absolute path set using the `--log-bin` option. An absolute path remains unchanged; in such a case, the index must be edited manually to enable the new path or paths to be used. (In older versions of MySQL, manual intervention was required whenever relocating the binary log or relay log files.) (Bug #11745230, Bug #12133)

Setting this option causes the `log_bin` system variable to be set to `ON` (or `1`), and not to the base name. The binary log file name (with path) is available as the `log_bin_basename` system variable.

If you specify this option without also specifying a `--server-id`, the server is not allowed to start. (Bug #11763963, Bug #56739)

- `--log-bin-index[=file_name]`

Command-Line Format	`--log-bin-index=file_name`	
Permitted Values	**Type**	file name

The index file for binary log file names. See Section 5.4.4, “The Binary Log”. If you omit the file name, and if you did not specify one with `--log-bin`, MySQL uses *`host_name`*`-bin.index` as the file name.

- ~~`--log-bin-trust-function-creators[={0|1}]`~~

Command-Line Format	--log-bin-trust-function-creators	
System Variable	Name	log_bin_trust_function_creators
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	FALSE

This option sets the corresponding log_bin_trust_function_creators system variable. If no argument is given, the option sets the variable to 1. log_bin_trust_function_creators affects how MySQL enforces restrictions on stored function and trigger creation. See Section 23.7, “Binary Logging of Stored Programs”.

- --log-bin-use-v1-row-events[={0|1}]

Command-Line Format	--log-bin-use-v1-row-events[={0|1}]	
System Variable	Name	log_bin_use_v1_row_events
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	0

MySQL 8.0 uses Version 2 binary log row events, which cannot be read by MySQL Server releases prior to MySQL 5.6.6. Setting this option to 1 causes mysqld to write the binary log using Version 1 logging events, which is the only version of binary log events used in previous releases, and thus produce binary logs that can be read by older slaves. Setting --log-bin-use-v1-row-events to 0 (the default) causes mysqld to use Version 2 binary log events.

The value used for this option can be obtained from the read-only log_bin_use_v1_row_events system variable.

Statement selection options. The options in the following list affect which statements are written to the binary log, and thus sent by a replication master server to its slaves. There are also options for slave servers that control which statements received from the master should be executed or ignored. For details, see Section 18.1.6.3, “Replication Slave Options and Variables”.

- --binlog-do-db=*db_name*

Command-Line Format	--binlog-do-db=name	
Permitted Values	Type	string

This option affects binary logging in a manner similar to the way that --replicate-do-db affects replication.

The effects of this option depend on whether the statement-based or row-based logging format is in use, in the same way that the effects of --replicate-do-db depend on whether statement-based or row-based replication is in use. You should keep in mind that the format used to log a given statement may not necessarily be the same as that indicated by the value of binlog_format. For example, DDL statements such as CREATE TABLE and ALTER TABLE are always logged as statements, without regard to the logging format in effect, so the following statement-based rules for --binlog-do-db always apply in determining whether or not the statement is logged.

Statement-based logging. Only those statements are written to the binary log where the default database (that is, the one selected by `USE`) is *`db_name`*. To specify more than one database, use this option multiple times, once for each database; however, doing so does *not* cause cross-database statements such as `UPDATE some_db.some_table SET foo='bar'` to be logged while a different database (or no database) is selected.

Warning

To specify multiple databases you *must* use multiple instances of this option. Because database names can contain commas, the list will be treated as the name of a single database if you supply a comma-separated list.

An example of what does not work as you might expect when using statement-based logging: If the server is started with `--binlog-do-db=sales` and you issue the following statements, the `UPDATE` statement is *not* logged:

```
USE prices;
UPDATE sales.january SET amount=amount+1000;
```

The main reason for this “just check the default database” behavior is that it is difficult from the statement alone to know whether it should be replicated (for example, if you are using multiple-table `DELETE` statements or multiple-table `UPDATE` statements that act across multiple databases). It is also faster to check only the default database rather than all databases if there is no need.

Another case which may not be self-evident occurs when a given database is replicated even though it was not specified when setting the option. If the server is started with `--binlog-do-db=sales`, the following `UPDATE` statement is logged even though `prices` was not included when setting `--binlog-do-db`:

```
USE sales;
UPDATE prices.discounts SET percentage = percentage + 10;
```

Because `sales` is the default database when the `UPDATE` statement is issued, the `UPDATE` is logged.

Row-based logging. Logging is restricted to database *`db_name`*. Only changes to tables belonging to *`db_name`* are logged; the default database has no effect on this. Suppose that the server is started with `--binlog-do-db=sales` and row-based logging is in effect, and then the following statements are executed:

```
USE prices;
UPDATE sales.february SET amount=amount+100;
```

The changes to the `february` table in the `sales` database are logged in accordance with the `UPDATE` statement; this occurs whether or not the `USE` statement was issued. However, when using the row-based logging format and `--binlog-do-db=sales`, changes made by the following `UPDATE` are not logged:

```
USE prices;
UPDATE prices.march SET amount=amount-25;
```

Even if the `USE prices` statement were changed to `USE sales`, the `UPDATE` statement's effects would still not be written to the binary log.

Another important difference in `--binlog-do-db` handling for statement-based logging as opposed to the row-based logging occurs with regard to statements that refer to multiple databases. Suppose that the server is started with `--binlog-do-db=db1`, and the following statements are executed:

```
USE db1;
UPDATE db1.table1 SET col1 = 10, db2.table2 SET col2 = 20;
```

If you are using statement-based logging, the updates to both tables are written to the binary log. However, when using the row-based format, only the changes to `table1` are logged; `table2` is in a different database, so it is not changed by the `UPDATE`. Now suppose that, instead of the `USE db1` statement, a `USE db4` statement had been used:

```
USE db4;
UPDATE db1.table1 SET col1 = 10, db2.table2 SET col2 = 20;
```

In this case, the `UPDATE` statement is not written to the binary log when using statement-based logging. However, when using row-based logging, the change to `table1` is logged, but not that to `table2`—in other words, only changes to tables in the database named by `--binlog-do-db` are logged, and the choice of default database has no effect on this behavior.

- `--binlog-ignore-db=`*`db_name`*

Command-Line Format	`--binlog-ignore-db=name`	
Permitted Values	Type	string

This option affects binary logging in a manner similar to the way that `--replicate-ignore-db` affects replication.

The effects of this option depend on whether the statement-based or row-based logging format is in use, in the same way that the effects of `--replicate-ignore-db` depend on whether statement-based or row-based replication is in use. You should keep in mind that the format used to log a given statement may not necessarily be the same as that indicated by the value of `binlog_format`. For example, DDL statements such as `CREATE TABLE` and `ALTER TABLE` are always logged as statements, without regard to the logging format in effect, so the following statement-based rules for `--binlog-ignore-db` always apply in determining whether or not the statement is logged.

Statement-based logging. Tells the server to not log any statement where the default database (that is, the one selected by `USE`) is *`db_name`*.

When there is no default database, no `--binlog-ignore-db` options are applied, and such statements are always logged. (Bug #11829838, Bug #60188)

Row-based format. Tells the server not to log updates to any tables in the database *`db_name`*. The current database has no effect.

When using statement-based logging, the following example does not work as you might expect. Suppose that the server is started with `--binlog-ignore-db=sales` and you issue the following statements:

```
USE prices;
UPDATE sales.january SET amount=amount+1000;
```

The `UPDATE` statement *is* logged in such a case because `--binlog-ignore-db` applies only to the default database (determined by the `USE` statement). Because the `sales` database was specified explicitly in the statement, the statement has not been filtered. However, when using row-based logging, the `UPDATE` statement's effects are *not* written to the binary log, which means that no changes to the `sales.january` table are logged; in this instance, `--binlog-ignore-db=sales` causes *all* changes made to tables in the master's copy of the `sales` database to be ignored for purposes of binary logging.

To specify more than one database to ignore, use this option multiple times, once for each database. Because database names can contain commas, the list will be treated as the name of a single database if you supply a comma-separated list.

You should not use this option if you are using cross-database updates and you do not want these updates to be logged.

Checksum options. MySQL 8.0 supports reading and writing of binary log checksums. These are enabled using the two options listed here:

- `--binlog-checksum={NONE|CRC32}`

Command-Line Format	`--binlog-checksum=type`	
Permitted Values	Type	string
	Default	`CRC32`
	Valid Values	`NONE`
		`CRC32`

Enabling this option causes the master to write checksums for events written to the binary log. Set to `NONE` to disable, or the name of the algorithm to be used for generating checksums; currently, only CRC32 checksums are supported, and CRC32 is the default.

- `--master-verify-checksum={0|1}`

Command-Line Format	`--master-verify-checksum=name`	
Permitted Values	Type	boolean
	Default	`OFF`

Enabling this option causes the master to verify events from the binary log using checksums, and to stop with an error in the event of a mismatch. Disabled by default.

To control reading of checksums by the slave (from the relay) log, use the `--slave-sql-verify-checksum` option.

Testing and debugging options. The following binary log options are used in replication testing and debugging. They are not intended for use in normal operations.

- `--max-binlog-dump-events=N`

Command-Line Format	`--max-binlog-dump-events=#`	
Permitted Values	Type	integer
	Default	`0`

This option is used internally by the MySQL test suite for replication testing and debugging.

- `--sporadic-binlog-dump-fail`

Command-Line Format	`--sporadic-binlog-dump-fail`	
Permitted Values	Type	boolean
	Default	`FALSE`

This option is used internally by the MySQL test suite for replication testing and debugging.

- `--binlog-rows-query-log-events`

Command-Line Format	`--binlog-rows-query-log-events`

Permitted Values	Type	boolean
	Default	FALSE

This option enables binlog_rows_query_log_events.

System Variables Used with Binary Logging

The following list describes system variables for controlling binary logging. They can be set at server startup and some of them can be changed at runtime using SET. Server options used to control binary logging are listed earlier in this section. For information about the sql_log_bin and sql_log_off variables, see Section 5.1.5, “Server System Variables”.

- binlog_cache_size

Command-Line Format	--binlog-cache-size=#	
System Variable	Name	binlog_cache_size
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	32768
	Min Value	4096
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	32768
	Min Value	4096
	Max Value	18446744073709551615

The size of the cache to hold changes to the binary log during a transaction. A binary log cache is allocated for each client if the server supports any transactional storage engines and if the server has the binary log enabled (--log-bin option). If you often use large transactions, you can increase this cache size to get better performance. The Binlog_cache_use and Binlog_cache_disk_use status variables can be useful for tuning the size of this variable. See Section 5.4.4, “The Binary Log”.

binlog_cache_size sets the size for the transaction cache only; the size of the statement cache is governed by the binlog_stmt_cache_size system variable.

- binlog_checksum

System Variable	Name	binlog_checksum
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string
	Default	CRC32

	Valid Values	`NONE`
		`CRC32`

When enabled, this variable causes the master to write a checksum for each event in the binary log. `binlog_checksum` supports the values `NONE` (disabled) and `CRC32`. The default is `CRC32`.

When `binlog_checksum` is disabled (value `NONE`), the server verifies that it is writing only complete events to the binary log by writing and checking the event length (rather than a checksum) for each event.

Changing the value of this variable causes the binary log to be rotated; checksums are always written to an entire binary log file, and never to only part of one.

Setting this variable on the master to a value unrecognized by the slave causes the slave to set its own `binlog_checksum` value to `NONE`, and to stop replication with an error. (Bug #13553750, Bug #61096) If backward compatibility with older slaves is a concern, you may want to set the value explicitly to `NONE`.

- `binlog_direct_non_transactional_updates`

Command-Line Format	`--binlog-direct-non-transactional-updates[=value]`	
System Variable	**Name**	`binlog_direct_non_transactional_updates`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Due to concurrency issues, a slave can become inconsistent when a transaction contains updates to both transactional and nontransactional tables. MySQL tries to preserve causality among these statements by writing nontransactional statements to the transaction cache, which is flushed upon commit. However, problems arise when modifications done to nontransactional tables on behalf of a transaction become immediately visible to other connections because these changes may not be written immediately into the binary log.

The `binlog_direct_non_transactional_updates` variable offers one possible workaround to this issue. By default, this variable is disabled. Enabling `binlog_direct_non_transactional_updates` causes updates to nontransactional tables to be written directly to the binary log, rather than to the transaction cache.

`binlog_direct_non_transactional_updates` works only for statements that are replicated using the statement-based binary logging format; that is, it works only when the value of `binlog_format` is `STATEMENT`, or when `binlog_format` is `MIXED` and a given statement is being replicated using the statement-based format. This variable has no effect when the binary log format is `ROW`, or when `binlog_format` is set to `MIXED` and a given statement is replicated using the row-based format.

Important

Before enabling this variable, you must make certain that there are no dependencies between transactional and nontransactional tables; an example of such a dependency would be the statement `INSERT INTO myisam_table SELECT * FROM innodb_table`. Otherwise, such statements are likely to cause the slave to diverge from the master.

In MySQL 8.0, this variable has no effect when the binary log format is `ROW` or `MIXED`. (Bug #51291)

- `binlog_error_action`

Command-Line Format	`--binlog-error-action[=value]`	
System Variable	Name	`binlog_error_action`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`IGNORE_ERROR`
	Valid Values	`IGNORE_ERROR`
		`ABORT_SERVER`
Permitted Values	Type	enumeration
	Default	`ABORT_SERVER`
	Valid Values	`IGNORE_ERROR`
		`ABORT_SERVER`

Controls what happens when the server encounters an error such as not being able to write to, flush or synchronize the binary log, which can cause the master's log to become inconsistent and replication slaves to lose synchronization.

This variable defaults to `ABORT_SERVER`, which makes the server halt logging and shut down whenever it encounters such an error with the binary log. Upon server restart, all of the previously prepared and binary logged transactions are committed, while any transactions which were prepared but not binary logged due to the error are aborted.

When `binlog_error_action` is set to `IGNORE_ERROR`, if the server encounters such an error it continues the ongoing transaction, logs the error then halts logging, and continues performing updates. To resume binary logging `log_bin` must be enabled again. This provides backward compatibility with older versions of MySQL.

In previous releases this variable was named `binlogging_impossible_mode`.

- `binlog_format`

Command-Line Format	`--binlog-format=format`	
System Variable	Name	`binlog_format`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`ROW`
	Valid Values	`ROW`
		`STATEMENT`
		`MIXED`

This variable sets the binary logging format, and can be any one of `STATEMENT`, `ROW`, or `MIXED`. See Section 18.2.1, “Replication Formats”. `binlog_format` is set by the `--binlog-format` option at startup, or by the `binlog_format` variable at runtime.

Note

While you can change the logging format at runtime, it is *not* recommended that you change it while replication is ongoing. This is due in part to the fact that slaves do not honor the master's `binlog_format` setting; a given MySQL Server can change only its own logging format.

The default is `ROW`.

You must have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege to set either the global or session `binlog_format` value.

The rules governing when changes to this variable take effect and how long the effect lasts are the same as for other MySQL server system variables. For more information, see Section 13.7.4.1, "SET Syntax for Variable Assignment".

When `MIXED` is specified, statement-based replication is used, except for cases where only row-based replication is guaranteed to lead to proper results. For example, this happens when statements contain user-defined functions (UDF) or the `UUID()` function. An exception to this rule is that `MIXED` always uses statement-based replication for stored functions and triggers.

There are exceptions when you cannot switch the replication format at runtime:

- From within a stored function or a trigger.
- If the session is currently in row-based replication mode and has open temporary tables.
- From within a transaction.

Trying to switch the format in those cases results in an error.

The binary log format affects the behavior of the following server options:

- `--replicate-do-db`
- `--replicate-ignore-db`
- `--binlog-do-db`
- `--binlog-ignore-db`

These effects are discussed in detail in the descriptions of the individual options.

- `binlog_group_commit_sync_delay`

Command-Line Format	`--binlog-group-commit-sync-delay=#`	
System Variable	**Name**	`binlog_group_commit_sync_delay`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1000000`

Controls how many microseconds the binary log commit waits before synchronizing the binary log file to disk. By default `binlog-group-commit-sync-delay` is set to 0, meaning that there is no delay. Setting `binlog-group-commit-sync-delay` to a microsecond delay enables more transactions to be synchronized together to disk at once, reducing the overall time to commit a group of transactions because the larger groups require fewer time units per group. With the correct tuning, this can increase slave performance without compromising the master's throughput.

- `binlog_group_commit_sync_no_delay_count`

Command-Line Format	`--binlog-group-commit-sync-no-delay-count=#`	
System Variable	**Name**	`binlog_group_commit_sync_no_delay_count`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`1000000`

The maximum number of transactions to wait for before aborting the current delay as specified by `binlog-group-commit-sync-delay`. If `binlog-group-commit-sync-delay` is set to 0, then this option has no effect.

- `binlogging_impossible_mode`

Deprecated	5.7.6	
Command-Line Format	`--binlogging-impossible-mode[=value]`	
System Variable	**Name**	`binlogging_impossible_mode`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`IGNORE_ERROR`
	Valid Values	`IGNORE_ERROR`
		`ABORT_SERVER`

This option is deprecated and will be removed in a future MySQL release. Use the renamed `binlog_error_action` to control what happens when the server cannot write to the binary log.

- `binlog_max_flush_queue_time`

Deprecated	5.7.9	
System Variable	**Name**	`binlog_max_flush_queue_time`
	Variable Scope	Global
	Dynamic Variable	Yes

Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`100000`

Formerly, this controlled the time in microseconds to continue reading transactions from the flush queue before proceeding with group commit. In MySQL 8.0, this variable no longer has any effect.

`binlog_max_flush_queue_time` is deprecated as of MySQL 5.7.9, and is marked for eventual removal in a future MySQL release.

- `binlog_order_commits`

System Variable	**Name**	`binlog_order_commits`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

When this variable is enabled on a master (the default), transactions are externalized in the same order as they are written to the binary log. If disabled, transactions may be committed in parallel. In some cases, disabling this variable might produce a performance increment.

- `binlog_row_image`

Command-Line Format	`--binlog-row-image=image_type`	
System Variable	**Name**	`binlog_row_image=image_type`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`full`
	Valid Values	`full` (Log all columns)
		`minimal` (Log only changed columns, and columns needed to identify rows)
		`noblob` (Log all columns, except for unneeded BLOB and TEXT columns)

In MySQL row-based replication, each row change event contains two images, a “before” image whose columns are matched against when searching for the row to be updated, and an “after” image containing the changes. Normally, MySQL logs full rows (that is, all columns) for both the before and after images. However, it is not strictly necessary to include every column in both images, and we can often save disk, memory, and network usage by logging only those columns which are actually required.

Note

When deleting a row, only the before image is logged, since there are no changed values to propagate following the deletion. When inserting a row, only the after image is logged, since there is no existing row to be matched. Only when updating a row are both the before and after images required, and both written to the binary log.

For the before image, it is necessary only that the minimum set of columns required to uniquely identify rows is logged. If the table containing the row has a primary key, then only the primary key column or columns are written to the binary log. Otherwise, if the table has a unique key all of whose columns are `NOT NULL`, then only the columns in the unique key need be logged. (If the table has neither a primary key nor a unique key without any `NULL` columns, then all columns must be used in the before image, and logged.) In the after image, it is necessary to log only the columns which have actually changed.

You can cause the server to log full or minimal rows using the `binlog_row_image` system variable. This variable actually takes one of three possible values, as shown in the following list:

- `full`: Log all columns in both the before image and the after image.
- `minimal`: Log only those columns in the before image that are required to identify the row to be changed; log only those columns in the after image that are actually changed.
- `noblob`: Log all columns (same as `full`), except for `BLOB` and `TEXT` columns that are not required to identify rows, or that have not changed.

The default value is `full`.

When using `minimal` or `noblob`, deletes and updates are guaranteed to work correctly for a given table if and only if the following conditions are true for both the source and destination tables:

- All columns must be present and in the same order; each column must use the same data type as its counterpart in the other table.
- The tables must have identical primary key definitions.

(In other words, the tables must be identical with the possible exception of indexes that are not part of the tables' primary keys.)

If these conditions are not met, it is possible that the primary key column values in the destination table may prove insufficient to provide a unique match for a delete or update. In this event, no warning or error is issued; the master and slave silently diverge, thus breaking consistency.

Setting this variable has no effect when the binary logging format is `STATEMENT`. When `binlog_format` is `MIXED`, the setting for `binlog_row_image` is applied to changes that are logged using row-based format, but this setting no effect on changes logged as statements.

Setting `binlog_row_image` on either the global or session level does not cause an implicit commit; this means that this variable can be changed while a transaction is in progress without affecting the transaction.

- `binlog_row_metadata`

Introduced	8.0.1	
Command-Line Format	`--binlog-row-metadata=metadata_type`	
System Variable	**Name**	`binlog_row_metadata=metadata_type`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`MINIMAL`
	Valid Values	`FULL` (All metadata is included.)
		`MINIMAL` (Limit included metadata.)

Configures the amount of table metadata added to the binary log when using row-based logging. When set to `MINIMAL`, the default, only metadata related to `SIGNED` flags, column character set and geometry types are logged. When set to `FULL` complete metadata for tables is logged, such as column name, `ENUM` or `SET` string values, `PRIMARY KEY` information, and so on.

The extended metadata serves the following purposes:

- Slaves use the metadata to transfer data when its table structure is different from the master's.
- External software can use the metadata to decode row events and store the data into external databases, such as a data warehouse.

- `binlog_rows_query_log_events`

System Variable	**Name**	`binlog_rows_query_log_events`
	Variable Scope	Global, Session
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`FALSE`

The `binlog_rows_query_log_events` system variable affects row-based logging only. When enabled, it causes the MySQL Server to write informational log events such as row query log events into its binary log. This information can be used for debugging and related purposes; such as obtaining the original query issued on the master when it cannot be reconstructed from the row updates.

These events are normally ignored by MySQL programs reading the binary log and so cause no issues when replicating or restoring from backup. To view them, increase the verbosity level by using mysqlbinlog's `--verbose` option twice, either as "-vv" or "--verbose --verbose".

- `binlog_stmt_cache_size`

Command-Line Format		`--binlog-stmt-cache-size=#`
System Variable	**Name**	`binlog_stmt_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`32768`
	Min Value	`4096`

	Max Value	4294967295
Permitted Values (64-bit platforms)	**Type**	integer
	Default	32768
	Min Value	4096
	Max Value	18446744073709551615

This variable determines the size of the cache for the binary log to hold nontransactional statements issued during a transaction. Separate binary log transaction and statement caches are allocated for each client if the server supports any transactional storage engines and if the server has the binary log enabled (`--log-bin` option). If you often use large nontransactional statements during transactions, you can increase this cache size to get better performance. The `Binlog_stmt_cache_use` and `Binlog_stmt_cache_disk_use` status variables can be useful for tuning the size of this variable. See Section 5.4.4, “The Binary Log”.

The `binlog_cache_size` system variable sets the size for the transaction cache.

- `binlog_transaction_dependency_tracking`

Introduced	8.0.1	
Command-Line Format	`--binlog-transaction-dependency-tracking=value`	
System Variable	**Name**	`binlog_transaction_dependency_tracking`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`COMMIT_ORDER`
	Valid Values	`WRITESET`
		`WRITESET_SESSION`

The source of dependency information on the master used to determine which transactions can be executed in parallel by the slave's multi-threaded applier. This variable can take one of the three values described in the following list:

- `COMMIT_ORDER`: Dependency information is generated from the master's commit timestamps. This is the default. This mode is also used for any transactions without write sets, even if this vairable's is `WRITESET` or `WRITESET_SESSION`; this is also the case for transactions updating tables without primary keys and transactions updating tables having foreign key constraints.
- `WRITESET`: Dependency information is generated from the master's write set, and any transactions which write different tuples can be parallelized.
- `WRITESET_SESSION`: Dependency information is generated from the master's write set, but no two updates from the same session can be reordered.

`WRITESET` and `WRITESET_SESSION` modes do not deliver any transaction dependencies that are newer then those that would have been returned in `COMMIT_ORDER` mode.

The value of this variable cannot be set to anything other than `COMMIT_ORDER` if `transaction_write_set_extraction` is `OFF`. You should also note that the value

of `transaction_write_set_extraction` cannot be changed if the current value of `binlog_transaction_dependency_tracking` is `WRITESET` or `WRITESET_SESSION`.

The number of row hashes to be kept and checked for the latest transaction to have changed a given row is determined by the value of `binlog_transaction_dependency_history_size`.

- `binlog_transaction_dependency_history_size`

Introduced	8.0.1	
Command-Line Format	`--binlog-transaction-dependency-size=#`	
System Variable	**Name**	`binlog_transaction_dependency_history_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`25000`
	Min Value	`1`
	Max Value	`1000000`

Sets an upper limit on the number of row hashes which are kept in memory and used for looking up the transaction that last modified a given row. Once this number of hashes has been reached, the history is purged.

- `log_bin`

System Variable	**Name**	`log_bin`
	Variable Scope	Global
	Dynamic Variable	No

Whether the binary log is enabled. If the `--log-bin` option is used, then the value of this variable is `ON`; otherwise it is `OFF`. This variable reports only on the status of binary logging (enabled or disabled); it does not actually report the value to which `--log-bin` is set.

See Section 5.4.4, “The Binary Log”.

- `log_bin_basename`

System Variable	**Name**	`log_bin_basename`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name
	Default	`datadir + '/' + hostname + '-bin'`

Holds the name and complete path to the binary log file. Unlike the `log_bin` system variable, `log_bin_basename` reflects the name set with the `--log-bin` server option.

- `log_bin_index`

System Variable	Name	log_bin_index
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	file name

The index file for binary log file names.

- log_bin_use_v1_row_events

Command-Line Format	--log-bin-use-v1-row-events[={0|1}]	
System Variable	Name	log_bin_use_v1_row_events
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	0

Shows whether Version 2 binary logging is in use. A value of 1 shows that the server is writing the binary log using Version 1 logging events (the only version of binary log events used in previous releases), and thus producing a binary log that can be read by older slaves. 0 indicates that Version 2 binary log events are in use.

This variable is read-only. To switch between Version 1 and Version 2 binary event binary logging, it is necessary to restart mysqld with the --log-bin-use-v1-row-events option.

- log_slave_updates

Command-Line Format	--log-slave-updates	
System Variable	Name	log_slave_updates
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	FALSE

Whether updates received by a slave server from a master server should be logged to the slave's own binary log. Binary logging must be enabled on the slave for this variable to have any effect. See Section 18.1.6, “Replication and Binary Logging Options and Variables”.

- log_statements_unsafe_for_binlog

System Variable	Name	log_statements_unsafe_for_binlog
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean

	Default	`ON`

If error 1592 is encountered, controls whether the generated warnings are added to the error log or not.

- `master_verify_checksum`

System Variable	**Name**	`master_verify_checksum`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Enabling this variable causes the master to examine checksums when reading from the binary log. `master_verify_checksum` is disabled by default; in this case, the master uses the event length from the binary log to verify events, so that only complete events are read from the binary log.

- `max_binlog_cache_size`

Command-Line Format	`--max-binlog-cache-size=#`	
System Variable	**Name**	`max_binlog_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`18446744073709551615`
	Min Value	`4096`
	Max Value	`18446744073709551615`

If a transaction requires more than this many bytes of memory, the server generates a `Multi-statement transaction required more than 'max_binlog_cache_size' bytes of storage` error. The minimum value is 4096. The maximum possible value is 16EB (exabytes). The maximum recommended value is 4GB; this is due to the fact that MySQL currently cannot work with binary log positions greater than 4GB.

`max_binlog_cache_size` sets the size for the transaction cache only; the upper limit for the statement cache is governed by the `max_binlog_stmt_cache_size` system variable.

In MySQL 8.0, the visibility to sessions of `max_binlog_cache_size` matches that of the `binlog_cache_size` system variable; in other words, changing its value effects only new sessions that are started after the value is changed.

- `max_binlog_size`

Command-Line Format	`--max-binlog-size=#`	
System Variable	**Name**	`max_binlog_size`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1073741824`
	Min Value	`4096`
	Max Value	`1073741824`

If a write to the binary log causes the current log file size to exceed the value of this variable, the server rotates the binary logs (closes the current file and opens the next one). The minimum value is 4096 bytes. The maximum and default value is 1GB.

A transaction is written in one chunk to the binary log, so it is never split between several binary logs. Therefore, if you have big transactions, you might see binary log files larger than `max_binlog_size`.

If `max_relay_log_size` is 0, the value of `max_binlog_size` applies to relay logs as well.

- `max_binlog_stmt_cache_size`

Command-Line Format	`--max-binlog-stmt-cache-size=#`	
System Variable	**Name**	`max_binlog_stmt_cache_size`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`18446744073709547520`
	Min Value	`4096`
	Max Value	`18446744073709547520`

If nontransactional statements within a transaction require more than this many bytes of memory, the server generates an error. The minimum value is 4096. The maximum and default values are 4GB on 32-bit platforms and 16EB (exabytes) on 64-bit platforms.

`max_binlog_stmt_cache_size` sets the size for the statement cache only; the upper limit for the transaction cache is governed exclusively by the `max_binlog_cache_size` system variable.

- `original_commit_timestamp`

Introduced	8.0.1	
Status Variable	**Name**	`original_commit_timestamp`
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	**Type**	numeric

For internal use by replication. When re-executing a transaction on a slave, this is set to the time when the transaction was committed on the original master, measured in microseconds since the epoch. This allows the original commit timestamp to be propagated throughout a replication topology.

- `sync_binlog`

Command-Line Format	`--sync-binlog=#`	
System Variable	**Name**	`sync_binlog`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (32-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`
Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

Controls the number of binary log commit groups to collect before synchronizing the binary log to disk. When `sync_binlog=0`, the binary log is never synchronized to disk, and when `sync_binlog` is set to a value greater than 0 this number of binary log commit groups is periodically synchronized to disk. When `sync_binlog=1`, all transactions are synchronized to the binary log before they are committed. Therefore, even in the event of an unexpected restart, any transactions that are missing from the binary log are only in prepared state. This causes the server's automatic recovery routine to roll back those transactions. This guarantees that no transaction is lost from the binary log, and is the safest option. However this can have a negative impact on performance because of an increased number of disk writes. Using a higher value improves performance, but with the increased risk of data loss.

When `sync_binlog=0` or `sync_binlog` is greater than 1, transactions are committed without having been synchronized to disk. Therefore, in the event of a power failure or operating system crash, it is possible that the server has committed some transactions that have not been synchronized to the binary log. Therefore it is impossible for the recovery routine to recover these transactions and they will be lost from the binary log.

The default value of `sync_binlog` is 1, which is the safest choice, but as noted above can impact performance.

Prior to MySQL 5.7.7, the default value was 0, which configures no synchronizing to disk—in this case, the server relies on the operating system to flush the binary log's contents from time to time as for any other file.

18.1.6.5 Global Transaction ID Options and Variables

Startup Options Used with GTID Replication

System Variables Used with GTID Replication

The MySQL Server options and system variables described in this section are used to monitor and control Global Transaction Identifiers (GTIDs).

For additional information, see Section 18.1.3, “Replication with Global Transaction Identifiers”.

Startup Options Used with GTID Replication

The following server startup options are used with GTID-based replication:

- `--enforce-gtid-consistency`

Command-Line Format	`--enforce-gtid-consistency[=value]`	
System Variable	Name	`enforce_gtid_consistency`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`OFF`
	Valid Values	`OFF`
		`ON`
		`WARN`

When enabled, the server enforces GTID consistency by allowing execution of only statements that can be safely logged using a GTID. You *must* set this option to `ON` before enabling GTID based replication.

The values that `--enforce-gtid-consistency` can be configured to are:

- `OFF`: all transactions are allowed to violate GTID consistency.
- `ON`: no transaction is allowed to violate GTID consistency.
- `WARN`: all transactions are allowed to violate GTID consistency, but a warning is generated in this case.

Setting `--enforce-gtid-consistency` without a value is an alias for `--enforce-gtid-consistency=ON`. This impacts on the behavior of the variable, see `enforce_gtid_consistency`.

Only statements that can be logged using GTID safe statements can be logged when `enforce-gtid-consistency` is set to `ON`, so the operations listed here cannot be used with this option:

- `CREATE TABLE ... SELECT` statements
- `CREATE TEMPORARY TABLE` or `DROP TEMPORARY TABLE` statements inside transactions

- Transactions or statements that update both transactional and nontransactional tables. There is an exception that nontransactional DML is allowed in the same transaction or in the same statement as transactional DML, if all *nontransactional* tables are temporary.

For more information, see Section 18.1.3.4, “Restrictions on Replication with GTIDs”.

- `--executed-gtids-compression-period`

Deprecated	5.7.6	
Command-Line Format	`--executed-gtids-compression-period=#`	
Permitted Values	**Type**	integer
	Default	`1000`
	Min Value	`0`
	Max Value	`4294967295`

This option is deprecated and will be removed in a future MySQL release. Use the renamed gtid_executed_compression_period to control how the gtid_executed table is compressed.

- `--gtid-mode`

Command-Line Format	`--gtid-mode=MODE`	
System Variable	**Name**	`gtid_mode`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`OFF`
	Valid Values	`OFF`
		`OFF_PERMISSIVE`
		`ON_PERMISSIVE`
		`ON`

This option specifies whether global transaction identifiers (GTIDs) are used to identify transactions. Setting this option to `--gtid-mode=ON` requires that `enforce-gtid-consistency` be set to `ON`. The `gtid_mode` variable is dynamic and enables GTID based replication to be configured online. Before using this feature, see Section 18.1.5, “Changing Replication Modes on Online Servers”.

- `--gtid-executed-compression-period`

Command-Line Format	`--gtid-executed-compression-period=#`	
Permitted Values	**Type**	integer
	Default	`1000`
	Min Value	`0`
	Max Value	`4294967295`

Compress the `mysql.gtid_executed` table each time this many transactions have taken place. A setting of 0 means that this table is not compressed. No compression of the table occurs when binary logging is enabled, therefore the option has no effect unless `log_bin` is `OFF`.

See mysql.gtid_executed Table Compression, for more information.

System Variables Used with GTID Replication

The following system variables are used with GTID-based replication:

- `binlog_gtid_simple_recovery`

Command-Line Format		`--binlog-gtid-simple-recovery`
System Variable	Name	`binlog_gtid_simple_recovery`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	`FALSE`
Permitted Values	Type	boolean
	Default	`TRUE`

This variable controls how binary log files are iterated during the search for GTIDs when MySQL starts or restarts.

When `binlog_gtid_simple_recovery=FALSE`, the method of iterating the binary log files is:

- To initialize `gtid_executed`, binary log files are iterated from the newest file, stopping at the first binary log that has any `Previous_gtids_log_event`. All GTIDs from `Previous_gtids_log_event` and `Gtid_log_events` are read from this binary log file. This GTID set is stored internally and called `gtids_in_binlog`. The value of `gtid_executed` is computed as the union of this set and the GTIDs stored in the `mysql.gtid_executed` table.

  This process could take a long time if you had a large number of binary log files without GTID events, for example created when `gtid_mode=OFF`.

- To initialize `gtid_purged`, binary log files are iterated from the oldest to the newest, stopping at the first binary log that contains either a `Previous_gtids_log_event` that is non-empty (that has at least one GTID) or that has at least one `Gtid_log_event`. From this binary log it reads `Previous_gtids_log_event`. This GTID set is subtracted from `gtids_in_binlog` and the result stored in the internal variable `gtids_in_binlog_not_purged`. The value of `gtid_purged` is initialized to the value of `gtid_executed`, minus `gtids_in_binlog_not_purged`.

When `binlog_gtid_simple_recovery=TRUE`, which is the default, the server iterates only the oldest and the newest binary log files and the values of `gtid_purged` and `gtid_executed` are computed based only on `Previous_gtids_log_event` or `Gtid_log_event` found in these files. This ensures only two binary log files are iterated during server restart or when binary logs are being purged.

Note

If this option is enabled, `gtid_executed` and `gtid_purged` may be initialized incorrectly in the following situations:

> - The newest binary log was generated by MySQL 5.7.5 or older, and `gtid_mode` was `ON` for some binary logs but `OFF` for the newest binary log.
> - A `SET GTID_PURGED` statement was issued on a MySQL version prior to 5.7.7, and the binary log that was active at the time of the `SET GTID_PURGED` has not yet been purged.
>
> If an incorrect GTID set is computed in either situation, it will remain incorrect even if the server is later restarted, regardless of the value of this option.

- `enforce_gtid_consistency`

Command-Line Format	`--enforce-gtid-consistency[=value]`	
System Variable	Name	`enforce_gtid_consistency`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`OFF`
	Valid Values	`OFF`
		`ON`
		`WARN`

Depending on the value of this variable, the server enforces GTID consistency by allowing execution of only statements that can be safely logged using a GTID. You *must* set this variable to `ON` before enabling GTID based replication.

The values that `enforce_gtid_consistency` can be configured to are:

- `OFF`: all transactions are allowed to violate GTID consistency.
- `ON`: no transaction is allowed to violate GTID consistency.
- `WARN`: all transactions are allowed to violate GTID consistency, but a warning is generated in this case.

For more information on statements that can be logged using GTID based replication, see `--enforce-gtid-consistency`.

Prior to MySQL 5.7.6, the boolean `enforce-gtid-consistency` defaulted to `OFF`. To maintain compatibility with previous versions, in MySQL 5.7.6 the enumeration defaults to `OFF`, and setting `--enforce-gtid-consistency` without a value is interpreted as setting the value to `ON`. The variable also has multiple textual aliases for the values: `0=OFF=FALSE`, `1=ON=TRUE`,`2=WARN`. This differs from other enumeration types but maintains compatibility with the boolean type used in previous versions. These changes impact on what is returned by the variable. Using `SELECT @@ENFORCE_GTID_CONSISTENCY`, `SHOW VARIABLES LIKE 'ENFORCE_GTID_CONSISTENCY'`, and `SELECT * FROM INFORMATION_SCHEMA.VARIABLES WHERE 'VARIABLE_NAME' = 'ENFORCE_GTID_CONSISTENCY'`, all return the textual form, not the numeric form. This is an incompatible change, since `@@ENFORCE_GTID_CONSISTENCY` returns the numeric form for booleans but returns the textual form for `SHOW` and the Information Schema.

- `executed_gtids_compression_period`

Deprecated	5.7.6

System Variable	Name	executed_gtids_compression_period
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	1000
	Min Value	0
	Max Value	4294967295

This option is deprecated and will be removed in a future MySQL release. Use the renamed gtid_executed_compression_period to control how the gtid_executed table is compressed.

- gtid_executed

System Variable	Name	gtid_executed
	Variable Scope	Global, Session
	Dynamic Variable	No
System Variable	Name	gtid_executed
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	string

When used with global scope, this variable contains a representation of the set of all transactions executed on the server and GTIDs that have been set by a SET gtid_purged statement. This is the same as the value of the Executed_Gtid_Set column in the output of SHOW MASTER STATUS and SHOW SLAVE STATUS. The value of this variable is a GTID set, see GTID Sets for more information.

When the server starts, @@global.gtid_executed is initialized. See binlog_gtid_simple_recovery for more information on how binary logs are iterated to populate gtid_executed. GTIDs are then added to the set as transactions are executed, or if any SET gtid_purged statement is executed.

The set of transactions that can be found in the binary logs at any given time is equal to GTID_SUBTRACT(@@global.gtid_executed, @@global.gtid_purged); that is, to all transactions in the binary log that have not yet been purged.

Issuing RESET MASTER causes the global value (but not the session value) of this variable to be reset to an empty string. GTIDs are not otherwise removed from this set other than when the set is cleared due to RESET MASTER.

Prior to MySQL 5.7.7, this variable could also be used with session scope, where it contained a representation of the set of transactions that are written to the cache in the current session. The session scope was deprecated in MySQL 5.7.7.

- gtid_executed_compression_period

System Variable	Name	gtid_executed_compression_period
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	1000
	Min Value	0
	Max Value	4294967295

Compress the mysql.gtid_executed table each time this many transactions have been processed. A setting of 0 means that this table is not compressed. Since no compression of the table occurs when using the binary log, setting the value of the variable has no effect unless binary logging is disabled.

See mysql.gtid_executed Table Compression, for more information.

- gtid_mode

System Variable	Name	gtid_mode
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	OFF
	Valid Values	OFF
		OFF_PERMISSIVE
		ON_PERMISSIVE
		ON

Controls whether GTID based logging is enabled and what type of transactions the logs can contain. Prior to MySQL 5.7.6 this variable was read-only and was set using the --gtid-mode option only. MySQL 5.7.6 enables this variable to be set dynamically. You must have the SYSTEM_VARIABLES_ADMIN or SUPER privilege to set this variable. enforce_gtid_consistency must be true before you can set gtid_mode=ON. Before modifying this variable, see Section 18.1.5, “Changing Replication Modes on Online Servers”.

Logged transactions can be either anonymous or use GTIDs. Anonymous transactions rely on binary log file and position to identify specific transactions. GTID transactions have a unique identifier that is used to refer to transactions. The different modes are:

- OFF: Both new and replicated transactions must be anonymous.
- OFF_PERMISSIVE: New transactions are anonymous. Replicated transactions can be either anonymous or GTID transactions.
- ON_PERMISSIVE: New transactions are GTID transactions. Replicated transactions can be either anonymous or GTID transactions.
- ON: Both new and replicated transactions must be GTID transactions.

Changes from one value to another can only be one step at a time. For example, if `gtid_mode` is currently set to `OFF_PERMISSIVE`, it is possible to change to `OFF` or `ON_PERMISSIVE` but not to `ON`.

The values of `gtid_purged` and `gtid_executed` are persistent regardless of the value of `gtid_mode`. Therefore even after changing the value of `gtid_mode`, these variables contain the correct values.

- `gtid_next`

System Variable	Name	gtid_next
	Variable Scope	Session
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	AUTOMATIC
	Valid Values	AUTOMATIC
		ANONYMOUS
		UUID:NUMBER

This variable is used to specify whether and how the next GTID is obtained. `gtid_next` can take any of the following values:

- `AUTOMATIC`: Use the next automatically-generated global transaction ID.
- `ANONYMOUS`: Transactions do not have global identifiers, and are identified by file and position only.
- A global transaction ID in *`UUID`*:*`NUMBER`* format.

Exactly which of the above options are valid depends on the setting of `gtid_mode`, see Section 18.1.5.1, “Replication Mode Concepts” for more information. Setting this variable has no effect if `gtid_mode` is `OFF`.

After this variable has been set to *`UUID`*:*`NUMBER`*, and a transaction has been committed or rolled back, an explicit `SET GTID_NEXT` statement must again be issued before any other statement.

`DROP TABLE` or `DROP TEMPORARY TABLE` fails with an explicit error when used on a combination of nontemporary tables with temporary tables, or of temporary tables using transactional storage engines with temporary tables using nontransactional storage engines.

- `gtid_owned`

System Variable	Name	gtid_owned
	Variable Scope	Global, Session
	Dynamic Variable	No
Permitted Values	Type	string

This read-only variable holds a list whose contents depend on its scope. When used with session scope, the list holds all GTIDs that are owned by this client; when used with global scope, it holds a list of all GTIDs along with their owners.

System Variable	Name	gtid_purged
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

The set of all transactions that have been purged from the binary log. This is a subset of the set of transactions in gtid_executed. The value of this variable is a GTID set, see GTID Sets for more information.

When the server starts, the global value of gtid_purged is initialized to a set of GTIDs. See binlog_gtid_simple_recovery for more information on how binary logs are iterated to populate gtid_purged. Issuing RESET MASTER causes the value of this variable to be reset to an empty string.

There are two ways to set gtid_purged. When gtid_set is a superset of gtid_purged, and goes not intersect with gtid_subtract(gtid_executed - gtid_purged) use:

```
SET @@GLOBAL.GTID_PURGED = 'gtid_set'
```

The result is that GTID_PURGED is set equal to gtid_set, and GTID_EXECUTED becomes the union of gtid_set and the previous value of GTID_EXECUTED.

When gtid_set does not intersect with gtid_executed use:

```
SET @@GLOBAL.GTID_PURGED = '+gtid_set'
```

The result is that gtid_set is added to both gtid_executed and gtid_purged.

If all existing binary logs were generated using MySQL 5.7.6 or later, after issuing a SET gtid_purged statement, binlog_gtid_simple_recovery=TRUE (the default setting in MySQL 5.7.7 and later) can safely be used. If binary logs from MySQL 5.7.7 or earlier exist, there is a chance that gtid_purged may be computed incorrectly. See binlog_gtid_simple_recovery for more information. If you are using MySQL 5.7.7 or earlier, after issuing a SET gtid_purged statement note down the current binary log file name, which can be checked using SHOW MASTER STATUS. If the server is restarted before this file has been purged, then you should use binlog_gtid_simple_recovery=FALSE to avoid gtid_purged or gtid_executed being computed incorrectly.

- simplified_binlog_gtid_recovery

Deprecated	5.7.6	
Command-Line Format	--simplified-binlog-gtid-recovery	
System Variable	Name	simplified_binlog_gtid_recovery
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	boolean
	Default	FALSE

This option is deprecated and will be removed in a future MySQL release. Use the renamed `binlog_gtid_simple_recovery` to control how MySQL iterates through binary log files after a crash.

18.1.7 Common Replication Administration Tasks

Once replication has been started it executes without requiring much regular administration. This section describes how to check the status of replication and how to pause a slave.

18.1.7.1 Checking Replication Status

The most common task when managing a replication process is to ensure that replication is taking place and that there have been no errors between the slave and the master. The primary statement for this is `SHOW SLAVE STATUS`, which you must execute on each slave:

```
mysql> SHOW SLAVE STATUS\G
*************************** 1. row ***************************
               Slave_IO_State: Waiting for master to send event
                  Master_Host: master1
                  Master_User: root
                  Master_Port: 3306
                Connect_Retry: 60
              Master_Log_File: mysql-bin.000004
          Read_Master_Log_Pos: 931
               Relay_Log_File: slave1-relay-bin.000056
                Relay_Log_Pos: 950
        Relay_Master_Log_File: mysql-bin.000004
             Slave_IO_Running: Yes
            Slave_SQL_Running: Yes
              Replicate_Do_DB:
          Replicate_Ignore_DB:
           Replicate_Do_Table:
       Replicate_Ignore_Table:
      Replicate_Wild_Do_Table:
  Replicate_Wild_Ignore_Table:
                   Last_Errno: 0
                   Last_Error:
                 Skip_Counter: 0
          Exec_Master_Log_Pos: 931
              Relay_Log_Space: 1365
              Until_Condition: None
               Until_Log_File:
                Until_Log_Pos: 0
           Master_SSL_Allowed: No
           Master_SSL_CA_File:
           Master_SSL_CA_Path:
              Master_SSL_Cert:
            Master_SSL_Cipher:
               Master_SSL_Key:
        Seconds_Behind_Master: 0
Master_SSL_Verify_Server_Cert: No
                Last_IO_Errno: 0
                Last_IO_Error:
               Last_SQL_Errno: 0
               Last_SQL_Error:
  Replicate_Ignore_Server_Ids: 0
```

The key fields from the status report to examine are:

- `Slave_IO_State`: The current status of the slave. See Section 8.14.5, “Replication Slave I/O Thread States”, and Section 8.14.6, “Replication Slave SQL Thread States”, for more information.
- `Slave_IO_Running`: Whether the I/O thread for reading the master's binary log is running. Normally, you want this to be `Yes` unless you have not yet started replication or have explicitly stopped it with `STOP SLAVE`.

- `Slave_SQL_Running`: Whether the SQL thread for executing events in the relay log is running. As with the I/O thread, this should normally be `Yes`.
- `Last_IO_Error`, `Last_SQL_Error`: The last errors registered by the I/O and SQL threads when processing the relay log. Ideally these should be blank, indicating no errors.
- `Seconds_Behind_Master`: The number of seconds that the slave SQL thread is behind processing the master binary log. A high number (or an increasing one) can indicate that the slave is unable to handle events from the master in a timely fashion.

  A value of 0 for `Seconds_Behind_Master` can usually be interpreted as meaning that the slave has caught up with the master, but there are some cases where this is not strictly true. For example, this can occur if the network connection between master and slave is broken but the slave I/O thread has not yet noticed this—that is, `slave_net_timeout` has not yet elapsed.

  It is also possible that transient values for `Seconds_Behind_Master` may not reflect the situation accurately. When the slave SQL thread has caught up on I/O, `Seconds_Behind_Master` displays 0; but when the slave I/O thread is still queuing up a new event, `Seconds_Behind_Master` may show a large value until the SQL thread finishes executing the new event. This is especially likely when the events have old timestamps; in such cases, if you execute `SHOW SLAVE STATUS` several times in a relatively short period, you may see this value change back and forth repeatedly between 0 and a relatively large value.

Several pairs of fields provide information about the progress of the slave in reading events from the master binary log and processing them in the relay log:

- (`Master_Log_file`, `Read_Master_Log_Pos`): Coordinates in the master binary log indicating how far the slave I/O thread has read events from that log.
- (`Relay_Master_Log_File`, `Exec_Master_Log_Pos`): Coordinates in the master binary log indicating how far the slave SQL thread has executed events received from that log.
- (`Relay_Log_File`, `Relay_Log_Pos`): Coordinates in the slave relay log indicating how far the slave SQL thread has executed the relay log. These correspond to the preceding coordinates, but are expressed in slave relay log coordinates rather than master binary log coordinates.

On the master, you can check the status of connected slaves using `SHOW PROCESSLIST` to examine the list of running processes. Slave connections have `Binlog Dump` in the `Command` field:

```
mysql> SHOW PROCESSLIST \G;
*************************** 4. row ***************************
     Id: 10
   User: root
   Host: slave1:58371
     db: NULL
Command: Binlog Dump
   Time: 777
  State: Has sent all binlog to slave; waiting for binlog to be updated
   Info: NULL
```

Because it is the slave that drives the replication process, very little information is available in this report.

For slaves that were started with the `--report-host` option and are connected to the master, the `SHOW SLAVE HOSTS` statement on the master shows basic information about the slaves. The output includes the ID of the slave server, the value of the `--report-host` option, the connecting port, and master ID:

```
mysql> SHOW SLAVE HOSTS;
+-----------+--------+------+------------------+-----------+
| Server_id | Host   | Port | Rpl_recovery_rank | Master_id |
+-----------+--------+------+------------------+-----------+
```

```
|        10 | slave1 | 3306 |                  0 |         1 |
+-----------+--------+------+--------------------+-----------+
1 row in set (0.00 sec)
```

18.1.7.2 Pausing Replication on the Slave

You can stop and start replication on the slave using the `STOP SLAVE` and `START SLAVE` statements.

To stop processing of the binary log from the master, use `STOP SLAVE`:

```
mysql> STOP SLAVE;
```

When replication is stopped, the slave I/O thread stops reading events from the master binary log and writing them to the relay log, and the SQL thread stops reading events from the relay log and executing them. You can pause the I/O or SQL thread individually by specifying the thread type:

```
mysql> STOP SLAVE IO_THREAD;
mysql> STOP SLAVE SQL_THREAD;
```

To start execution again, use the `START SLAVE` statement:

```
mysql> START SLAVE;
```

To start a particular thread, specify the thread type:

```
mysql> START SLAVE IO_THREAD;
mysql> START SLAVE SQL_THREAD;
```

For a slave that performs updates only by processing events from the master, stopping only the SQL thread can be useful if you want to perform a backup or other task. The I/O thread will continue to read events from the master but they are not executed. This makes it easier for the slave to catch up when you restart the SQL thread.

Stopping only the I/O thread enables the events in the relay log to be executed by the SQL thread up to the point where the relay log ends. This can be useful when you want to pause execution to catch up with events already received from the master, when you want to perform administration on the slave but also ensure that it has processed all updates to a specific point. This method can also be used to pause event receipt on the slave while you conduct administration on the master. Stopping the I/O thread but permitting the SQL thread to run helps ensure that there is not a massive backlog of events to be executed when replication is started again.

18.2 Replication Implementation

Replication is based on the master server keeping track of all changes to its databases (updates, deletes, and so on) in its binary log. The binary log serves as a written record of all events that modify database structure or content (data) from the moment the server was started. Typically, `SELECT` statements are not recorded because they modify neither database structure nor content.

Each slave that connects to the master requests a copy of the binary log. That is, it pulls the data from the master, rather than the master pushing the data to the slave. The slave also executes the events from the binary log that it receives. This has the effect of repeating the original changes just as they were made on the master. Tables are created or their structure modified, and data is inserted, deleted, and updated according to the changes that were originally made on the master.

Because each slave is independent, the replaying of the changes from the master's binary log occurs independently on each slave that is connected to the master. In addition, because each slave receives a copy of the binary log only by requesting it from the master, the slave is able to read and update the copy of the database at its own pace and can start and stop the replication process at will without affecting the ability to update to the latest database status on either the master or slave side.

For more information on the specifics of the replication implementation, see Section 18.2.2, “Replication Implementation Details”.

Masters and slaves report their status in respect of the replication process regularly so that you can monitor them. See Section 8.14, “Examining Thread Information”, for descriptions of all replicated-related states.

The master binary log is written to a local relay log on the slave before it is processed. The slave also records information about the current position with the master's binary log and the local relay log. See Section 18.2.4, “Replication Relay and Status Logs”.

Database changes are filtered on the slave according to a set of rules that are applied according to the various configuration options and variables that control event evaluation. For details on how these rules are applied, see Section 18.2.5, “How Servers Evaluate Replication Filtering Rules”.

18.2.1 Replication Formats

Replication works because events written to the binary log are read from the master and then processed on the slave. The events are recorded within the binary log in different formats according to the type of event. The different replication formats used correspond to the binary logging format used when the events were recorded in the master's binary log. The correlation between binary logging formats and the terms used during replication are:

- When using statement-based binary logging, the master writes SQL statements to the binary log. Replication of the master to the slave works by executing the SQL statements on the slave. This is called *statement-based replication* (often abbreviated as *SBR*), which corresponds to the standard MySQL statement-based binary logging format. Replication capabilities in MySQL version 5.1.4 and earlier used this format exclusively.

- When using row-based logging, the master writes *events* to the binary log that indicate how individual table rows are changed. Replication of the master to the slave works by copying the events representing the changes to the table rows to the slave. This is called *row-based replication* (often abbreviated as *RBR*).

  Row-based logging is the default method.

- You can also configure MySQL to use a mix of both statement-based and row-based logging, depending on which is most appropriate for the change to be logged. This is called *mixed-format logging*. When using mixed-format logging, a statement-based log is used by default. Depending on certain statements, and also the storage engine being used, the log is automatically switched to row-based in particular cases. Replication using the mixed format is often referred to as *mixed-based replication* or *mixed-format replication*. For more information, see Section 5.4.4.3, “Mixed Binary Logging Format”.

When using `MIXED` format, the binary logging format is determined in part by the storage engine being used and the statement being executed. For more information on mixed-format logging and the rules governing the support of different logging formats, see Section 5.4.4.3, “Mixed Binary Logging Format”.

The logging format in a running MySQL server is controlled by setting the `binlog_format` server system variable. This variable can be set with session or global scope. The rules governing when and how the new setting takes effect are the same as for other MySQL server system variables—setting the variable for the current session lasts only until the end of that session, and the change is not visible to other sessions; setting the variable globally requires a restart of the server to take effect. For more information, see Section 13.7.4.1, “SET Syntax for Variable Assignment”.

There are conditions under which you cannot change the binary logging format at runtime or doing so causes replication to fail. See Section 5.4.4.2, “Setting The Binary Log Format”.

You must have the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege to set either the global or session `binlog_format` value.

The statement-based and row-based replication formats have different issues and limitations. For a comparison of their relative advantages and disadvantages, see Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

With statement-based replication, you may encounter issues with replicating stored routines or triggers. You can avoid these issues by using row-based replication instead. For more information, see Section 23.7, “Binary Logging of Stored Programs”.

18.2.1.1 Advantages and Disadvantages of Statement-Based and Row-Based Replication

Each binary logging format has advantages and disadvantages. For most users, the mixed replication format should provide the best combination of data integrity and performance. If, however, you want to take advantage of the features specific to the statement-based or row-based replication format when performing certain tasks, you can use the information in this section, which provides a summary of their relative advantages and disadvantages, to determine which is best for your needs.

- Advantages of statement-based replication
- Disadvantages of statement-based replication
- Advantages of row-based replication
- Disadvantages of row-based replication

Advantages of statement-based replication

- Proven technology.
- Less data written to log files. When updates or deletes affect many rows, this results in *much* less storage space required for log files. This also means that taking and restoring from backups can be accomplished more quickly.
- Log files contain all statements that made any changes, so they can be used to audit the database.

Disadvantages of statement-based replication

- **Statements that are unsafe for SBR.**
  Not all statements which modify data (such as `INSERT DELETE`, `UPDATE`, and `REPLACE` statements) can be replicated using statement-based replication. Any nondeterministic behavior is difficult to replicate when using statement-based replication. Examples of such Data Modification Language (DML) statements include the following:
  - A statement that depends on a UDF or stored program that is nondeterministic, since the value returned by such a UDF or stored program or depends on factors other than the parameters supplied to it. (Row-based replication, however, simply replicates the value returned by the UDF or stored program, so its effect on table rows and data is the same on both the master and slave.) See Section 18.4.1.12, “Replication of Invoked Features”, for more information.
  - `DELETE` and `UPDATE` statements that use a `LIMIT` clause without an `ORDER BY` are nondeterministic. See Section 18.4.1.17, “Replication and LIMIT”.
  - Locking read statements (`SELECT ... FOR UPDATE` and `SELECT ... FOR SHARE`) that use `NOWAIT` or `SKIP LOCKED` options. See Locking Read Concurrency with NOWAIT and SKIP LOCKED.
  - Deterministic UDFs must be applied on the slaves.
  - Statements using any of the following functions cannot be replicated properly using statement-based replication:
    - `LOAD_FILE()`
    - `UUID()`, `UUID_SHORT()`

- `USER()`
- `FOUND_ROWS()`
- `SYSDATE()` (unless both the master and the slave are started with the `--sysdate-is-now` option)
- `GET_LOCK()`
- `IS_FREE_LOCK()`
- `IS_USED_LOCK()`
- `MASTER_POS_WAIT()`
- `RAND()`
- `RELEASE_LOCK()`
- `SLEEP()`
- `VERSION()`

  However, all other functions are replicated correctly using statement-based replication, including `NOW()` and so forth.

  For more information, see Section 18.4.1.16, “Replication and System Functions”.

Statements that cannot be replicated correctly using statement-based replication are logged with a warning like the one shown here:

```
[Warning] Statement is not safe to log in statement format.
```

A similar warning is also issued to the client in such cases. The client can display it using `SHOW WARNINGS`.

- `INSERT ... SELECT` requires a greater number of row-level locks than with row-based replication.
- `UPDATE` statements that require a table scan (because no index is used in the `WHERE` clause) must lock a greater number of rows than with row-based replication.
- For `InnoDB`: An `INSERT` statement that uses `AUTO_INCREMENT` blocks other nonconflicting `INSERT` statements.
- For complex statements, the statement must be evaluated and executed on the slave before the rows are updated or inserted. With row-based replication, the slave only has to modify the affected rows, not execute the full statement.
- If there is an error in evaluation on the slave, particularly when executing complex statements, statement-based replication may slowly increase the margin of error across the affected rows over time. See Section 18.4.1.28, “Slave Errors During Replication”.
- Stored functions execute with the same `NOW()` value as the calling statement. However, this is not true of stored procedures.
- Deterministic UDFs must be applied on the slaves.
- Table definitions must be (nearly) identical on master and slave. See Section 18.4.1.10, “Replication with Differing Table Definitions on Master and Slave”, for more information.

Advantages of row-based replication

- All changes can be replicated. This is the safest form of replication.

Note

Statements that update the information in the `mysql` database—such as `GRANT`, `REVOKE` and the manipulation of triggers, stored routines (including stored procedures), and views—are all replicated to slaves using statement-based replication.

For statements such as `CREATE TABLE ... SELECT`, a `CREATE` statement is generated from the table definition and replicated using statement-based format, while the row insertions are replicated using row-based format.

- Fewer row locks are required on the master, which thus achieves higher concurrency, for the following types of statements:
  - `INSERT ... SELECT`
  - `INSERT` statements with `AUTO_INCREMENT`
  - `UPDATE` or `DELETE` statements with `WHERE` clauses that do not use keys or do not change most of the examined rows.
- Fewer row locks are required on the slave for any `INSERT`, `UPDATE`, or `DELETE` statement.

Disadvantages of row-based replication

- RBR can generate more data that must be logged. To replicate a DML statement (such as an `UPDATE` or `DELETE` statement), statement-based replication writes only the statement to the binary log. By contrast, row-based replication writes each changed row to the binary log. If the statement changes many rows, row-based replication may write significantly more data to the binary log; this is true even for statements that are rolled back. This also means that making and restoring a backup can require more time. In addition, the binary log is locked for a longer time to write the data, which may cause concurrency problems. Use `binlog_row_image=minimal` to reduce the disadvantage considerably.
- Deterministic UDFs that generate large `BLOB` values take longer to replicate with row-based replication than with statement-based replication. This is because the `BLOB` column value is logged, rather than the statement generating the data.
- You cannot see on the slave what statements were received from the master and executed. However, you can see what data was changed using `mysqlbinlog` with the options `--base64-output=DECODE-ROWS` and `--verbose`.

  Alternatively, use the `binlog_rows_query_log_events` variable, which if enabled adds a `Rows_query` event with the statement to `mysqlbinlog` output when the `-vv` option is used.
- For tables using the `MyISAM` storage engine, a stronger lock is required on the slave for `INSERT` statements when applying them as row-based events to the binary log than when applying them as statements. This means that concurrent inserts on `MyISAM` tables are not supported when using row-based replication.

18.2.1.2 Usage of Row-Based Logging and Replication

MySQL uses statement-based logging (SBL), row-based logging (RBL) or mixed-format logging. The type of binary log used impacts the size and efficiency of logging.Therefore the choice between row-based replication (RBR) or statement-based replication (SBR) depends on your application and environment. This section describes known issues when using a row-based format log, and describes some best practices using it in replication.

For additional information, see Section 18.2.1, “Replication Formats”, and Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

- **Row-based logging of temporary tables.** As noted in Section 18.4.1.24, “Replication and Temporary Tables”, temporary tables are not replicated when using row-based format. When using mixed format logging, “safe” statements involving temporary tables are logged using statement-based format. For more information, see Section 18.2.1.1, “Advantages and Disadvantages of Statement-Based and Row-Based Replication”.

  Temporary tables are not replicated when using row-based format because there is no need. In addition, because temporary tables can be read only from the thread which created them, there is seldom if ever any benefit obtained from replicating them, even when using statement-based format.

  In MySQL 8.0, you can switch from statement-based to row-based binary logging mode even when temporary tables have been created. However, while using the row-based format, the MySQL server cannot determine the logging mode that was in effect when a given temporary table was created. For this reason, the server in such cases logs a `DROP TEMPORARY TABLE IF EXISTS` statement for each temporary table that still exists for a given client session when that session ends. While this means that it is possible that an unnecessary `DROP TEMPORARY TABLE` statement might be logged in some cases, the statement is harmless, and does not cause an error even if the table does not exist, due to the presence of the `IF EXISTS` option.

  Nontransactional DML statements involving temporary tables are allowed when using `binlog_format=ROW`, as long as any nontransactional tables affected by the statements are temporary tables (Bug #14272672).

- **RBL and synchronization of nontransactional tables.** When many rows are affected, the set of changes is split into several events; when the statement commits, all of these events are written to the binary log. When executing on the slave, a table lock is taken on all tables involved, and then the rows are applied in batch mode. Depending on the engine used for the slave's copy of the table, this may or may not be effective.

- **Latency and binary log size.** RBL writes changes for each row to the binary log and so its size can increase quite rapidly. This can significantly increase the time required to make changes on the slave that match those on the master. You should be aware of the potential for this delay in your applications.

- **Reading the binary log.** `mysqlbinlog` displays row-based events in the binary log using the `BINLOG` statement (see Section 13.7.6.1, “BINLOG Syntax”). This statement displays an event as a base 64-encoded string, the meaning of which is not evident. When invoked with the `--base64-output=DECODE-ROWS` and `--verbose` options, `mysqlbinlog` formats the contents of the binary log to be human readable. When binary log events were written in row-based format and you want to read or recover from a replication or database failure you can use this command to read contents of the binary log. For more information, see Section 4.6.8.2, “mysqlbinlog Row Event Display”.

- **Binary log execution errors and slave_exec_mode.** If `slave_exec_mode` is `IDEMPOTENT`, a failure to apply changes from RBL because the original row cannot be found does not trigger an error or cause replication to fail. This means that it is possible that updates are not applied on the slave, so that the master and slave are no longer synchronized. Latency issues and use of nontransactional tables with RBR when `slave_exec_mode` is `IDEMPOTENT` can cause the master and slave to diverge even further. For more information about `slave_exec_mode`, see Section 5.1.5, “Server System Variables”.

  Note

  `slave_exec_mode=IDEMPOTENT` is generally useful only for circular replication or multi-master replication with MySQL Cluster, for which `IDEMPOTENT` is the default value.

  For other scenarios, setting `slave_exec_mode` to `STRICT` is normally sufficient; this is the default value for storage engines other than `NDB`.

The NDBCLUSTER storage engine is currently not supported in MySQL 8.0. See MySQL NDB Cluster 7.5 and NDB Cluster 7.6.

- **Lack of binary log checksums.** RBL does not use checksums, so network, disk, and other errors may not be identified when processing the binary log. To ensure that data is transmitted without network corruption use SSL for replication connections. The CHANGE MASTER TO statement has options to enable replication over SSL. See also Section 13.4.2.1, "CHANGE MASTER TO Syntax", for general information about setting up MySQL with SSL.

- **Filtering based on server ID not supported.** In MySQL 8.0, you can filter based on server ID by using the IGNORE_SERVER_IDS option for the CHANGE MASTER TO statement. This option works with statement-based and row-based logging formats. Another method to filter out changes on some slaves is to use a WHERE clause that includes the relation @@server_id <> *id_value* clause with UPDATE and DELETE statements. For example, WHERE @@server_id <> 1. However, this does not work correctly with row-based logging. To use the server_id system variable for statement filtering, use statement-based logging.

- **Database-level replication options.** The effects of the --replicate-do-db, --replicate-ignore-db, and --replicate-rewrite-db options differ considerably depending on whether row-based or statement-based logging is used. Therefore, it is recommended to avoid database-level options and instead use table-level options such as --replicate-do-table and --replicate-ignore-table. For more information about these options and the impact replication format has on how they operate, see Section 18.1.6, "Replication and Binary Logging Options and Variables".

- **RBL, nontransactional tables, and stopped slaves.** When using row-based logging, if the slave server is stopped while a slave thread is updating a nontransactional table, the slave database can reach an inconsistent state. For this reason, it is recommended that you use a transactional storage engine such as InnoDB for all tables replicated using the row-based format. Use of STOP SLAVE or STOP SLAVE SQL_THREAD prior to shutting down the slave MySQL server helps prevent issues from occurring, and is always recommended regardless of the logging format or storage engine you use.

18.2.1.3 Determination of Safe and Unsafe Statements in Binary Logging

The "safeness" of a statement in MySQL Replication, refers to whether the statement and its effects can be replicated correctly using statement-based format. If this is true of the statement, we refer to the statement as *safe*; otherwise, we refer to it as *unsafe*.

In general, a statement is safe if it deterministic, and unsafe if it is not. However, certain nondeterministic functions are *not* considered unsafe (see Nondeterministic functions not considered unsafe, later in this section). In addition, statements using results from floating-point math functions—which are hardware-dependent—are always considered unsafe (see Section 18.4.1.13, "Replication and Floating-Point Values").

Handling of safe and unsafe statements. A statement is treated differently depending on whether the statement is considered safe, and with respect to the binary logging format (that is, the current value of binlog_format).

- When using row-based logging, no distinction is made in the treatment of safe and unsafe statements.

- When using mixed-format logging, statements flagged as unsafe are logged using the row-based format; statements regarded as safe are logged using the statement-based format.

- When using statement-based logging, statements flagged as being unsafe generate a warning to this effect. Safe statements are logged normally.

Each statement flagged as unsafe generates a warning. Formerly, if a large number of such statements were executed on the master, this could lead to excessively large error log files. To prevent this, MySQL 5.7 provides a warning suppression mechanism, which behaves as follows: Whenever the 50

most recent `ER_BINLOG_UNSAFE_STATEMENT` warnings have been generated more than 50 times in any 50-second period, warning suppression is enabled. When activated, this causes such warnings not to be written to the error log; instead, for each 50 warnings of this type, a note `The last warning was repeated N times in last S seconds` is written to the error log. This continues as long as the 50 most recent such warnings were issued in 50 seconds or less; once the rate has decreased below this threshold, the warnings are once again logged normally. Warning suppression has no effect on how the safety of statements for statement-based logging is determined, nor on how warnings are sent to the client. MySQL clients still receive one warning for each such statement.

For more information, see Section 18.2.1, “Replication Formats”.

Statements considered unsafe.
Statements with the following characteristics are considered unsafe:

- **Statements containing system functions that may return a different value on slave.** These functions include `FOUND_ROWS()`, `GET_LOCK()`, `IS_FREE_LOCK()`, `IS_USED_LOCK()`, `LOAD_FILE()`, `MASTER_POS_WAIT()`, `PASSWORD()`, `RAND()`, `RELEASE_LOCK()`, `ROW_COUNT()`, `SESSION_USER()`, `SLEEP()`, `SYSDATE()`, `SYSTEM_USER()`, `USER()`, `UUID()`, and `UUID_SHORT()`.

  Nondeterministic functions not considered unsafe. Although these functions are not deterministic, they are treated as safe for purposes of logging and replication: `CONNECTION_ID()`, `CURDATE()`, `CURRENT_DATE()`, `CURRENT_TIME()`, `CURRENT_TIMESTAMP()`, `CURTIME()`,, `LAST_INSERT_ID()`, `LOCALTIME()`, `LOCALTIMESTAMP()`, `NOW()`, `UNIX_TIMESTAMP()`, `UTC_DATE()`, `UTC_TIME()`, and `UTC_TIMESTAMP()`.

  For more information, see Section 18.4.1.16, “Replication and System Functions”.

- **References to system variables.** Most system variables are not replicated correctly using the statement-based format. See Section 18.4.1.38, “Replication and Variables”. For exceptions, see Section 5.4.4.3, “Mixed Binary Logging Format”.

- **UDFs.** Since we have no control over what a UDF does, we must assume that it is executing unsafe statements.

- **Fulltext plugin.** This plugin may behave differently on different MySQL servers; therefore, statements depending on it could have different results. For this reason, all statements relying on the fulltext plugin are treated as unsafe in MySQL.

- **Trigger or stored program updates a table having an AUTO_INCREMENT column.** This is unsafe because the order in which the rows are updated may differ on the master and the slave.

  In addition, an `INSERT` into a table that has a composite primary key containing an `AUTO_INCREMENT` column that is not the first column of this composite key is unsafe.

  For more information, see Section 18.4.1.1, “Replication and AUTO_INCREMENT”.

- **INSERT ... ON DUPLICATE KEY UPDATE statements on tables with multiple primary or unique keys.** When executed against a table that contains more than one primary or unique key, this statement is considered unsafe, being sensitive to the order in which the storage engine checks the keys, which is not deterministic, and on which the choice of rows updated by the MySQL Server depends.

  An `INSERT ... ON DUPLICATE KEY UPDATE` statement against a table having more than one unique or primary key is marked as unsafe for statement-based replication. (Bug #11765650, Bug #58637)

- **Updates using LIMIT.** The order in which rows are retrieved is not specified, and is therefore considered unsafe. See Section 18.4.1.17, “Replication and LIMIT”.

- **Accesses or references log tables.** The contents of the system log table may differ between master and slave.

- **Nontransactional operations after transactional operations.** Within a transaction, allowing any nontransactional reads or writes to execute after any transactional reads or writes is considered unsafe.

  For more information, see Section 18.4.1.33, “Replication and Transactions”.

- **Accesses or references self-logging tables.** All reads and writes to self-logging tables are considered unsafe. Within a transaction, any statement following a read or write to self-logging tables is also considered unsafe.

- **LOAD DATA INFILE statements.** `LOAD DATA INFILE` is treated as unsafe and when `binlog_format=mixed` the statement is logged in row-based format. When `binlog_format=statement` `LOAD DATA INFILE` does not generate a warning, unlike other unsafe statements.

For additional information, see Section 18.4.1, “Replication Features and Issues”.

18.2.2 Replication Implementation Details

MySQL replication capabilities are implemented using three threads, one on the master server and two on the slave:

- **Binlog dump thread.** The master creates a thread to send the binary log contents to a slave when the slave connects. This thread can be identified in the output of `SHOW PROCESSLIST` on the master as the `Binlog Dump` thread.

  The binary log dump thread acquires a lock on the master's binary log for reading each event that is to be sent to the slave. As soon as the event has been read, the lock is released, even before the event is sent to the slave.

- **Slave I/O thread.** When a `START SLAVE` statement is issued on a slave server, the slave creates an I/O thread, which connects to the master and asks it to send the updates recorded in its binary logs.

  The slave I/O thread reads the updates that the master's `Binlog Dump` thread sends (see previous item) and copies them to local files that comprise the slave's relay log.

  The state of this thread is shown as `Slave_IO_running` in the output of `SHOW SLAVE STATUS` or as `Slave_running` in the output of `SHOW STATUS`.

- **Slave SQL thread.** The slave creates an SQL thread to read the relay log that is written by the slave I/O thread and execute the events contained therein.

In the preceding description, there are three threads per master/slave connection. A master that has multiple slaves creates one binary log dump thread for each currently connected slave, and each slave has its own I/O and SQL threads.

A slave uses two threads to separate reading updates from the master and executing them into independent tasks. Thus, the task of reading statements is not slowed down if statement execution is slow. For example, if the slave server has not been running for a while, its I/O thread can quickly fetch all the binary log contents from the master when the slave starts, even if the SQL thread lags far behind. If the slave stops before the SQL thread has executed all the fetched statements, the I/O thread has at least fetched everything so that a safe copy of the statements is stored locally in the slave's relay logs, ready for execution the next time that the slave starts.

The `SHOW PROCESSLIST` statement provides information that tells you what is happening on the master and on the slave regarding replication. For information on master states, see Section 8.14.4, “Replication Master Thread States”. For slave states, see Section 8.14.5, “Replication Slave I/O Thread States”, and Section 8.14.6, “Replication Slave SQL Thread States”.

The following example illustrates how the three threads show up in the output from `SHOW PROCESSLIST`.

On the master server, the output from `SHOW PROCESSLIST` looks like this:

```
mysql> SHOW PROCESSLIST\G
*************************** 1. row ***************************
     Id: 2
   User: root
   Host: localhost:32931
     db: NULL
Command: Binlog Dump
   Time: 94
  State: Has sent all binlog to slave; waiting for binlog to
         be updated
   Info: NULL
```

Here, thread 2 is a `Binlog Dump` replication thread that services a connected slave. The `State` information indicates that all outstanding updates have been sent to the slave and that the master is waiting for more updates to occur. If you see no `Binlog Dump` threads on a master server, this means that replication is not running; that is, no slaves are currently connected.

On a slave server, the output from `SHOW PROCESSLIST` looks like this:

```
mysql> SHOW PROCESSLIST\G
*************************** 1. row ***************************
     Id: 10
   User: system user
   Host:
     db: NULL
Command: Connect
   Time: 11
  State: Waiting for master to send event
   Info: NULL
*************************** 2. row ***************************
     Id: 11
   User: system user
   Host:
     db: NULL
Command: Connect
   Time: 11
  State: Has read all relay log; waiting for the slave I/O
         thread to update it
   Info: NULL
```

The `State` information indicates that thread 10 is the I/O thread that is communicating with the master server, and thread 11 is the SQL thread that is processing the updates stored in the relay logs. At the time that `SHOW PROCESSLIST` was run, both threads were idle, waiting for further updates.

The value in the `Time` column can show how late the slave is compared to the master. See Section A.13, “MySQL 8.0 FAQ: Replication”. If sufficient time elapses on the master side without activity on the `Binlog Dump` thread, the master determines that the slave is no longer connected. As for any other client connection, the timeouts for this depend on the values of `net_write_timeout` and `net_retry_count`; for more information about these, see Section 5.1.5, “Server System Variables”.

The `SHOW SLAVE STATUS` statement provides additional information about replication processing on a slave server. See Section 18.1.7.1, “Checking Replication Status”.

18.2.3 Replication Channels

Replication channels represent the path of transactions flowing from a master to a slave. This section describes how channels can be used in a replication topology, and the impact they have on single-source replication.

To provide compatibity with previous versions, the MySQL server automatically creates on startup a default channel whose name is the empty string (`""`). This channel is always present; it cannot be created or destroyed by the user. If no other channels (having nonempty names) have been

created, replication statements act on the default channel only, so that all replication statements from older slaves function as expected (see Section 18.2.3.2, “Compatibility with Previous Replication Statements”. Statements applying to replication channels as described in this section can be used only when there is at least one named channel.

A replication channel encompasses the path of transactions transmitted from a master to a slave. In multi-source replication a slave opens multiple channels, one per master, and each channel has its own relay log and applier (SQL) threads. Once transactions are received by a replication channel's receiver (I/O) thread, they are added to the channel's relay log file and passed through to an applier thread. This enables channels to function independently.

A replication channel is also associated with a host name and port. You can assign multiple channels to the same combination of host name and port; in MySQL 8.0, the maximum number of channels that can be added to one slave in a multi-source replication topology is 256. Each replication channel must have a unique (nonempty) name (see Section 18.2.3.4, “Replication Channel Naming Conventions”). Channels can be configured independently.

18.2.3.1 Commands for Operations on a Single Channel

To enable existing MySQL replication statements to act on individual replication channels use the `FOR CHANNEL channel_name` clause with the following replication statements:

- `CHANGE MASTER TO`
- `START SLAVE`
- `STOP SLAVE`
- `SHOW RELAYLOG EVENTS`
- `FLUSH RELAY LOGS`
- `SHOW SLAVE STATUS`
- `RESET SLAVE`

Similarly, an additional `channel_name` parameter is introduced for the following functions:

- `MASTER_POS_WAIT()`
- `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()`

The following statements are disallowed for the `group_replication_recovery` channel.

- `START SLAVE`
- `STOP SLAVE`

18.2.3.2 Compatibility with Previous Replication Statements

When a replication slave has multiple channels and a `FOR CHANNEL channel_name` option is not specified, a valid statement generally acts on all available channels.

For example, the following statements behave as expected:

- `START SLAVE` starts replication threads for all channels, except the `group_replication_recovery` channel.
- `STOP SLAVE` stops replication threads for all the channels, except the `group_replication_recovery` channel.
- `SHOW SLAVE STATUS` reports the status for all channels.
- `FLUSH RELAY LOGS` flushes the relay logs for all channels.
- `RESET SLAVE` resets all channels.

Warning

Use `RESET SLAVE` with caution as this statement deletes all existing channels, purges their relay log files, and recreates only the default channel.

Some replication statements cannot operate on all channels. In this case, error 1964 `Multiple channels exist on the slave. Please provide channel name as an argument.` is generated. The following statements and functions generate this error when used in a multi-source replication topology and a `FOR CHANNEL` *`channel_name`* option is not used to specify which channel to act on:

- `SHOW RELAYLOG EVENTS`
- `CHANGE MASTER TO`
- `MASTER_POS_WAIT()`
- `WAIT_UNTIL_SQL_THREAD_AFTER_GTIDS()`
- `WAIT_FOR_EXECUTED_GTID_SET()`

Note that a default channel always exists in a single source replication topology, where statements and functions behave as in previous versions of MySQL.

18.2.3.3 Startup Options and Replication Channels

This section describes startup options which are impacted by the addition of replication channels.

The following startup options *must* be configured correctly to use multi-source replication.

- `--relay-log-info-repository`

  This must be set to `TABLE`. If this option is set to `FILE`, attempting to add more sources to a slave fails with **`ER_SLAVE_NEW_CHANNEL_WRONG_REPOSITORY`**.

- `--master-info-repository`

  This must be set to `TABLE`. If this option is set to `FILE`, attempting to add more sources to a slave fails with **`ER_SLAVE_NEW_CHANNEL_WRONG_REPOSITORY`**.

The following startup options now affect *all* channels in a replication topology.

- `--log-slave-updates`

  All transactions received by the slave (even from multiple sources) are written in the binary log.

- `--relay-log-purge`

  When set, each channel purges its own relay log automatically.

- `--slave_transaction_retries`

  Applier threads of all channels retry transactions.

- `--skip-slave-start`

  No replication threads start on any channels.

- `--slave-skip-errors`

  Execution continues and errors are skipped for all channels.

The values set for the following startup options apply on each channel; since these are `mysqld` startup options, they are applied on every channel.

- `--max-relay-log-size=size`

  Maximum size of the individual relay log file for each channel; after reaching this limit, the file is rotated.

- `--relay-log-space-limit=size`

  Upper limit for the total size of all relay logs combined, for each individual channel. For *N* channels, the combined size of these logs is limited to `relay_log_space_limit * N`.

- `--slave-parallel-workers=value`

  Number of slave parallel workers per channel.

- `--slave-checkpoint-group`

  Waiting time by an I/O thread for each source.

- `--relay-log-index=filename`

  Base name for each channel's relay log index file. See Section 18.2.3.4, “Replication Channel Naming Conventions”.

- `--relay-log=filename`

  Denotes the base name of each channel's relay log file. See Section 18.2.3.4, “Replication Channel Naming Conventions”.

- `--slave_net-timeout=N`

  This value is set per channel, so that each channel waits for *N* seconds to check for a broken connection.

- `--slave-skip-counter=N`

  This value is set per channel, so that each channel skips *N* events from its master.

18.2.3.4 Replication Channel Naming Conventions

This section describes how naming conventions are impacted by replication channels.

Each replication channel has a unique name which is a string with a maximum length of 64 characters and is case insensitive. Because channel names are used in slave tables, the character set used for these is always UTF-8. Although you are generally free to use any name for channels, the following names are reserved:

- `group_replication_applier`
- `group_replication_recovery`

The name you choose for a replication channel also influences the file names used by a multi-source replication slave. The relay log files and index files for each channel are named *`base_name`*`-relay-bin-`*`channel_name`*`.0000x`, where *`base_name`* is generally a host name (if not specified using `--log-bin`) and *`channel_name`* is the name of the channel logged to this file.

18.2.4 Replication Relay and Status Logs

During replication, a slave server creates several logs that hold the binary log events relayed from the master to the slave, and to record information about the current status and location within the relay log. There are three types of logs used in the process, listed here:

- The *relay log* consists of the events read from the binary log of the master and written by the slave I/O thread. Events in the relay log are executed on the slave as part of the SQL thread.

- The *master info log* contains status and current configuration information for the slave's connection to the master. This log holds information on the master host name, login credentials, and coordinates indicating how far the slave has read from the master's binary log.

  This log can be written to the `mysql.slave_master_info` table instead of a file, by starting the slave with `--master-info-repository=TABLE`.

- The *relay log info log* holds status information about the execution point within the slave's relay log.

  This log can be written to the `mysql.slave_relay_log_info` table instead of a file by starting the slave with `--relay-log-info-repository=TABLE`.

Crash-safe replication. In order for replication to be crash-safe when using tables for logging status and relay information, these tables must use a transactional storage engine, such as `InnoDB`. In MySQL 8.0 these tables are created using `InnoDB`.

Therefore, in order to guarantee crash safety on the slave, you must run the slave with `--relay-log-recovery` enabled, in addition to setting `--relay-log-info-repository` to `TABLE`.

In MySQL 8.0, a warning is given when `mysqld` is unable to initialize the replication logging tables, but the slave is allowed to continue starting. This situation is most likely to occur when upgrading from a version of MySQL that does not support slave logging tables to one in which they are supported.

In MySQL 8.0, execution of any statement requiring a write lock on either or both of the `slave_master_info` and `slave_relay_log_info` tables is disallowed while replication is ongoing, while statements that perform only reads are permitted at any time.

Important

Do not attempt to update or insert rows in the `slave_master_info` or `slave_relay_log_info` table manually. Doing so can cause undefined behavior, and is not supported.

18.2.4.1 The Slave Relay Log

The relay log, like the binary log, consists of a set of numbered files containing events that describe database changes, and an index file that contains the names of all used relay log files.

The term “relay log file” generally denotes an individual numbered file containing database events. The term “relay log” collectively denotes the set of numbered relay log files plus the index file.

Relay log files have the same format as binary log files and can be read using `mysqlbinlog` (see Section 4.6.8, “`mysqlbinlog` — Utility for Processing Binary Log Files”).

By default, relay log file names have the form *`host_name`*`-relay-bin.`*`nnnnnn`* in the data directory, where *`host_name`* is the name of the slave server host and *`nnnnnn`* is a sequence number. Successive relay log files are created using successive sequence numbers, beginning with `000001`. The slave uses an index file to track the relay log files currently in use. The default relay log index file name is *`host_name`*`-relay-bin.index` in the data directory.

The default relay log file and relay log index file names can be overridden with, respectively, the `--relay-log` and `--relay-log-index` server options (see Section 18.1.6, “Replication and Binary Logging Options and Variables”).

If a slave uses the default host-based relay log file names, changing a slave's host name after replication has been set up can cause replication to fail with the errors `Failed to open the relay log` and `Could not find target log during relay log initialization`. This is a known issue (see Bug #2122). If you anticipate that a slave's host name might change in the future (for example, if networking is set up on the slave such that its host name can be modified using DHCP), you can avoid this issue entirely by using the `--relay-log` and `--relay-log-index` options to specify relay log file names explicitly when you initially set up the slave. This will make the names independent of server host name changes.

If you encounter the issue after replication has already begun, one way to work around it is to stop the slave server, prepend the contents of the old relay log index file to the new one, and then restart the slave. On a Unix system, this can be done as shown here:

```
shell> cat new_relay_log_name.index >> old_relay_log_name.index
shell> mv old_relay_log_name.index new_relay_log_name.index
```

A slave server creates a new relay log file under the following conditions:

- Each time the I/O thread starts.
- When the logs are flushed; for example, with `FLUSH LOGS` or `mysqladmin flush-logs`.
- When the size of the current relay log file becomes "too large," determined as follows:
  - If the value of `max_relay_log_size` is greater than 0, that is the maximum relay log file size.
  - If the value of `max_relay_log_size` is 0, `max_binlog_size` determines the maximum relay log file size.

The SQL thread automatically deletes each relay log file as soon as it has executed all events in the file and no longer needs it. There is no explicit mechanism for deleting relay logs because the SQL thread takes care of doing so. However, `FLUSH LOGS` rotates relay logs, which influences when the SQL thread deletes them.

18.2.4.2 Slave Status Logs

A replication slave server creates two logs. By default, these logs are files named `master.info` and `relay-log.info` and created in the data directory. The names and locations of these files can be changed by using the `--master-info-file` and `--relay-log-info-file` options, respectively. In MySQL 8.0, either or both of these logs can also be written to tables in the `mysql` database by starting the server with the appropriate option: use `--master-info-repository` to have the master info log written to the `mysql.slave_master_info` table, and use `--relay-log-info-repository` to have the relay log info log written to the `mysql.slave_relay_log_info` table. See Section 18.1.6, "Replication and Binary Logging Options and Variables".

The two status logs contain information like that shown in the output of the `SHOW SLAVE STATUS` statement, which is discussed in Section 13.4.2, "SQL Statements for Controlling Slave Servers". Because the status logs are stored on disk, they survive a slave server's shutdown. The next time the slave starts up, it reads the two logs to determine how far it has proceeded in reading binary logs from the master and in processing its own relay logs.

The master info log file or table should be protected because it contains the password for connecting to the master. See Section 6.1.2.3, "Passwords and Logging".

The slave I/O thread updates the master info log. The following table shows the correspondence between the lines in the `master.info` file, the columns in the `mysql.slave_master_info` table, and the columns displayed by `SHOW SLAVE STATUS`.

Line in `master.info` File	`slave_master_info` Table Column	`SHOW SLAVE STATUS` Column	Description
1	`Number_of_lines`	[None]	Number of lines in the file, or columns in the table
2	`Master_log_name`	`Master_Log_File`	The name of the master binary log currently

Line in `master.info` File	`slave_master_info` Table Column	`SHOW SLAVE STATUS` Column	Description
			being read from the master
3	`Master_log_pos`	`Read_Master_Log_Pos`	The current position within the master binary log that have been read from the master
4	`Host`	`Master_Host`	The host name of the master
5	`User_name`	`Master_User`	The user name used to connect to the master
6	`User_password`	Password (not shown by `SHOW SLAVE STATUS`)	The password used to connect to the master
7	`Port`	`Master_Port`	The network port used to connect to the master
8	`Connect_retry`	`Connect_Retry`	The period (in seconds) that the slave will wait before trying to reconnect to the master
9	`Enabled_ssl`	`Master_SSL_Allowed`	Indicates whether the server supports SSL connections
10	`Ssl_ca`	`Master_SSL_CA_File`	The file used for the Certificate Authority (CA) certificate
11	`Ssl_capath`	`Master_SSL_CA_Path`	The path to the Certificate Authority (CA) certificates
12	`Ssl_cert`	`Master_SSL_Cert`	The name of the SSL certificate file
13	`Ssl_cipher`	`Master_SSL_Cipher`	The list of possible ciphers used in the handshake for the SSL connection

Line in `master.info` File	`slave_master_info` Table Column	`SHOW SLAVE STATUS` Column	Description
14	`Ssl_key`	`Master_SSL_Key`	The name of the SSL key file
15	`Ssl_verify_server_cert`	`Master_SSL_Verify_Server_Cert`	Whether to verify the server certificate
16	`Heartbeat`	[None]	Interval between replication heartbeats, in seconds
17	`Bind`	`Master_Bind`	Which of the slave's network interfaces should be used for connecting to the master
18	`Ignored_server_ids`	`Replicate_Ignore_Server_Ids`	The list of server IDs to be ignored. Note that for `Ignored_server_ids` the list of server IDs is preceded by the total number of server IDs to ignore.
19	`Uuid`	`Master_UUID`	The master's unique ID
20	`Retry_count`	`Master_Retry_Count`	Maximum number of reconnection attempts permitted
21	`Ssl_crl`	[None]	Path to an ssl certificate revocation list file
22	`Ssl_crl_path`	[None]	Path to a directory containing ssl certificate revocation list files
23	`Enabled_auto_position`	`Auto_position`	If autopositioning is in use or not
24	`Channel_name`	`Channel_name`	The name of the replication channel

The slave SQL thread updates the relay log info log. In MySQL 8.0, the `relay-log.info` file includes a line count and a replication delay value. The following table shows the correspondence between the lines in the `relay-log.info` file, the columns in the `mysql.slave_relay_log_info` table, and the columns displayed by `SHOW SLAVE STATUS`.

Line in `relay-log.info`	`slave_relay_log_info` Table Column	`SHOW SLAVE STATUS` Column	Description
1	`Number_of_lines`	[None]	Number of lines in the file or columns in the table
2	`Relay_log_name`	`Relay_Log_File`	The name of the current relay log file
3	`Relay_log_pos`	`Relay_Log_Pos`	The current position within the relay log file; events up to this position have been executed on the slave database
4	`Master_log_name`	`Relay_Master_Log_File`	The name of the master binary log file from which the events in the relay log file were read
5	`Master_log_pos`	`Exec_Master_Log_Pos`	The equivalent position within the master's binary log file of events that have already been executed
6	`Sql_delay`	`SQL_Delay`	The number of seconds that the slave must lag the master
7	`Number_of_workers`	[None]	The number of slave worker threads for executing replication events (transactions) in parallel
8	`Id`	[None]	ID used for internal purposes; currently this is always 1
9	`Channel_name`	Channel_name	The name of the replication channel

In older versions of MySQL (prior to MySQL 5.6), the `relay-log.info` file does not include a line count or a delay value (and the `slave_relay_log_info` table is not available).

Line	Status Column	Description
1	`Relay_Log_File`	The name of the current relay log file
2	`Relay_Log_Pos`	The current position within the relay log file; events up to this position have been executed on the slave database
3	`Relay_Master_Log_File`	The name of the master binary log file from which the events in the relay log file were read
4	`Exec_Master_Log_Pos`	The equivalent position within the master's binary log file of events that have already been executed

Note

If you downgrade a slave server to a version older than MySQL 5.6, the older server does not read the `relay-log.info` file correctly. To address this, modify the file in a text editor by deleting the initial line containing the number of lines.

The contents of the `relay-log.info` file and the states shown by the `SHOW SLAVE STATUS` statement might not match if the `relay-log.info` file has not been flushed to disk. Ideally, you should only view `relay-log.info` on a slave that is offline (that is, `mysqld` is not running). For a running system, you can use `SHOW SLAVE STATUS`, or query the `slave_master_info` and `slave_relay_log_info` tables if you are writing the status logs to tables.

When you back up the slave's data, you should back up these two status logs, along with the relay log files. The status logs are needed to resume replication after you restore the data from the slave. If you lose the relay logs but still have the relay log info log, you can check it to determine how far the SQL thread has executed in the master binary logs. Then you can use `CHANGE MASTER TO` with the `MASTER_LOG_FILE` and `MASTER_LOG_POS` options to tell the slave to re-read the binary logs from that point. Of course, this requires that the binary logs still exist on the master.

18.2.5 How Servers Evaluate Replication Filtering Rules

If a master server does not write a statement to its binary log, the statement is not replicated. If the server does log the statement, the statement is sent to all slaves and each slave determines whether to execute it or ignore it.

On the master, you can control which databases to log changes for by using the `--binlog-do-db` and `--binlog-ignore-db` options to control binary logging. For a description of the rules that servers use in evaluating these options, see Section 18.2.5.1, “Evaluation of Database-Level Replication and Binary Logging Options”. You should not use these options to control which databases and tables are replicated. Instead, use filtering on the slave to control the events that are executed on the slave.

On the slave side, decisions about whether to execute or ignore statements received from the master are made according to the `--replicate-*` options that the slave was started with. (See Section 18.1.6, “Replication and Binary Logging Options and Variables”.) The filters governed by these options can also be set dynamically using the `CHANGE REPLICATION FILTER` statement. The rules governing such filters are the same whether they are created on startup using `--replicate-*` options or while the slave server is running by `CHANGE REPLICATION FILTER`.

In the simplest case, when there are no `--replicate-*` options, the slave executes all statements that it receives from the master. Otherwise, the result depends on the particular options given.

Database-level options (`--replicate-do-db`, `--replicate-ignore-db`) are checked first; see Section 18.2.5.1, “Evaluation of Database-Level Replication and Binary Logging Options”, for a description of this process. If no database-level options are used, option checking proceeds to any table-level options that may be in use (see Section 18.2.5.2, “Evaluation of Table-Level Replication

Options”, for a discussion of these). If one or more database-level options are used but none are matched, the statement is not replicated.

For statements affecting databases only (that is, `CREATE DATABASE`, `DROP DATABASE`, and `ALTER DATABASE`), database-level options always take precedence over any `--replicate-wild-do-table` options. In other words, for such statements, `--replicate-wild-do-table` options are checked if and only if there are no database-level options that apply. This is a change in behavior from previous versions of MySQL, where the statement `CREATE DATABASE dbx` was not replicated if the slave had been started with `--replicate-do-db=dbx --replicate-wild-do-table=db%.t1`. (Bug #46110)

To make it easier to determine what effect an option set will have, it is recommended that you avoid mixing “do” and “ignore” options, or wildcard and nonwildcard options.

If any `--replicate-rewrite-db` options were specified, they are applied before the `--replicate-*` filtering rules are tested.

Note

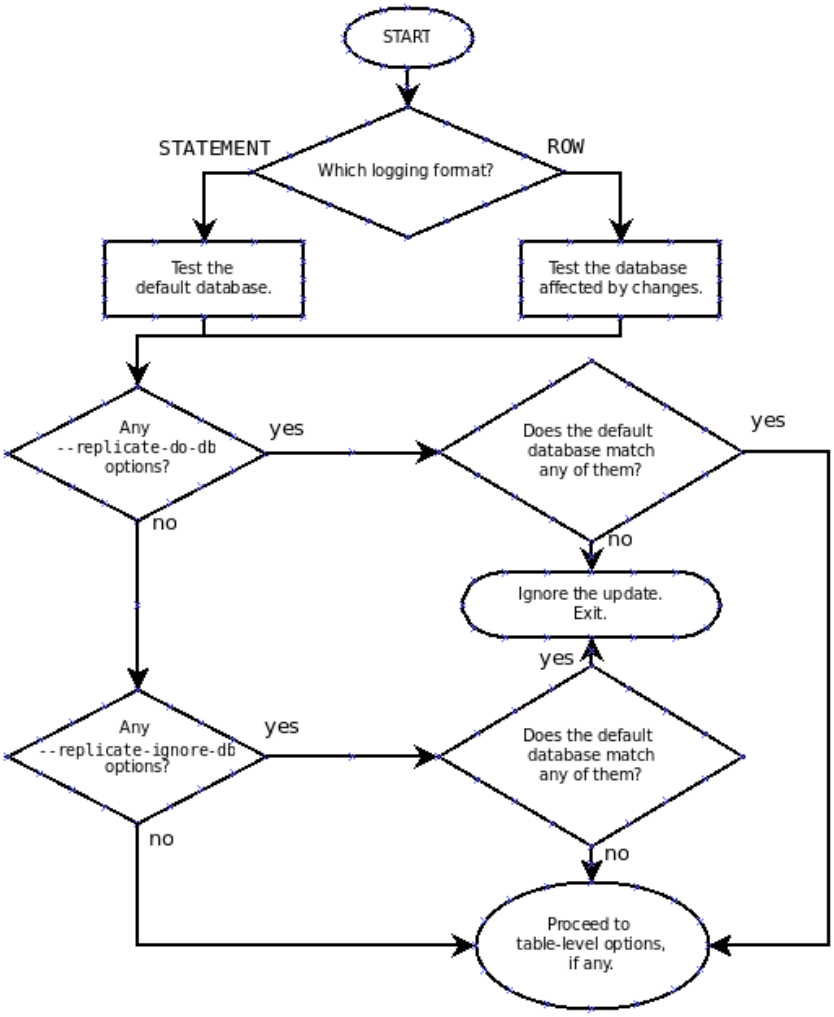
In MySQL 8.0, all replication filtering options follow the same rules for case sensitivity that apply to names of databases and tables elsewhere in the MySQL server, including the effects of the `lower_case_table_names` system variable.

This is a change from previous versions of MySQL. (Bug #51639)

18.2.5.1 Evaluation of Database-Level Replication and Binary Logging Options

When evaluating replication options, the slave begins by checking to see whether there are any `--replicate-do-db` or `--replicate-ignore-db` options that apply. When using `--binlog-do-db` or `--binlog-ignore-db`, the process is similar, but the options are checked on the master.

With statement-based replication, the default database is checked for a match. With row-based replication, the database where data is to be changed is the database that is checked. Regardless of the binary logging format, checking of database-level options proceeds as shown in the following diagram.



The steps involved are listed here:

1. Are there any `--replicate-do-db` options?
   - **Yes.** Do any of them match the database?
     - **Yes.** Execute the statement and exit.
     - **No.** Ignore the statement and exit.
   - **No.** Continue to step 2.
2. Are there any `--replicate-ignore-db` options?
   - **Yes.** Do any of them match the database?
     - **Yes.** Ignore the statement and exit.
     - **No.** Continue to step 3.
   - **No.** Continue to step 3.
3. Proceed to checking the table-level replication options, if there are any. For a description of how these options are checked, see Section 18.2.5.2, “Evaluation of Table-Level Replication Options”.

Important

A statement that is still permitted at this stage is not yet actually executed. The statement is not executed until all table-level options (if any) have also been checked, and the outcome of that process permits execution of the statement.

For binary logging, the steps involved are listed here:

1. Are there any `--binlog-do-db` or `--binlog-ignore-db` options?
   - **Yes.** Continue to step 2.
   - **No.** Log the statement and exit.
2. Is there a default database (has any database been selected by `USE`)?
   - **Yes.** Continue to step 3.
   - **No.** Ignore the statement and exit.
3. There is a default database. Are there any `--binlog-do-db` options?
   - **Yes.** Do any of them match the database?
     - **Yes.** Log the statement and exit.
     - **No.** Ignore the statement and exit.
   - **No.** Continue to step 4.
4. Do any of the `--binlog-ignore-db` options match the database?
   - **Yes.** Ignore the statement and exit.
   - **No.** Log the statement and exit.

Important

For statement-based logging, an exception is made in the rules just given for the `CREATE DATABASE`, `ALTER DATABASE`, and `DROP DATABASE` statements. In those cases, the database being *created, altered, or dropped* replaces the default database when determining whether to log or ignore updates.

`--binlog-do-db` can sometimes mean “ignore other databases”. For example, when using statement-based logging, a server running with only `--binlog-do-db=sales` does not write to the binary log statements for which the default database differs from `sales`. When using row-based logging with the same option, the server logs only those updates that change data in `sales`.

18.2.5.2 Evaluation of Table-Level Replication Options

The slave checks for and evaluates table options only if either of the following two conditions is true:

- No matching database options were found.
- One or more database options were found, and were evaluated to arrive at an “execute” condition according to the rules described in the previous section (see Section 18.2.5.1, “Evaluation of Database-Level Replication and Binary Logging Options”).

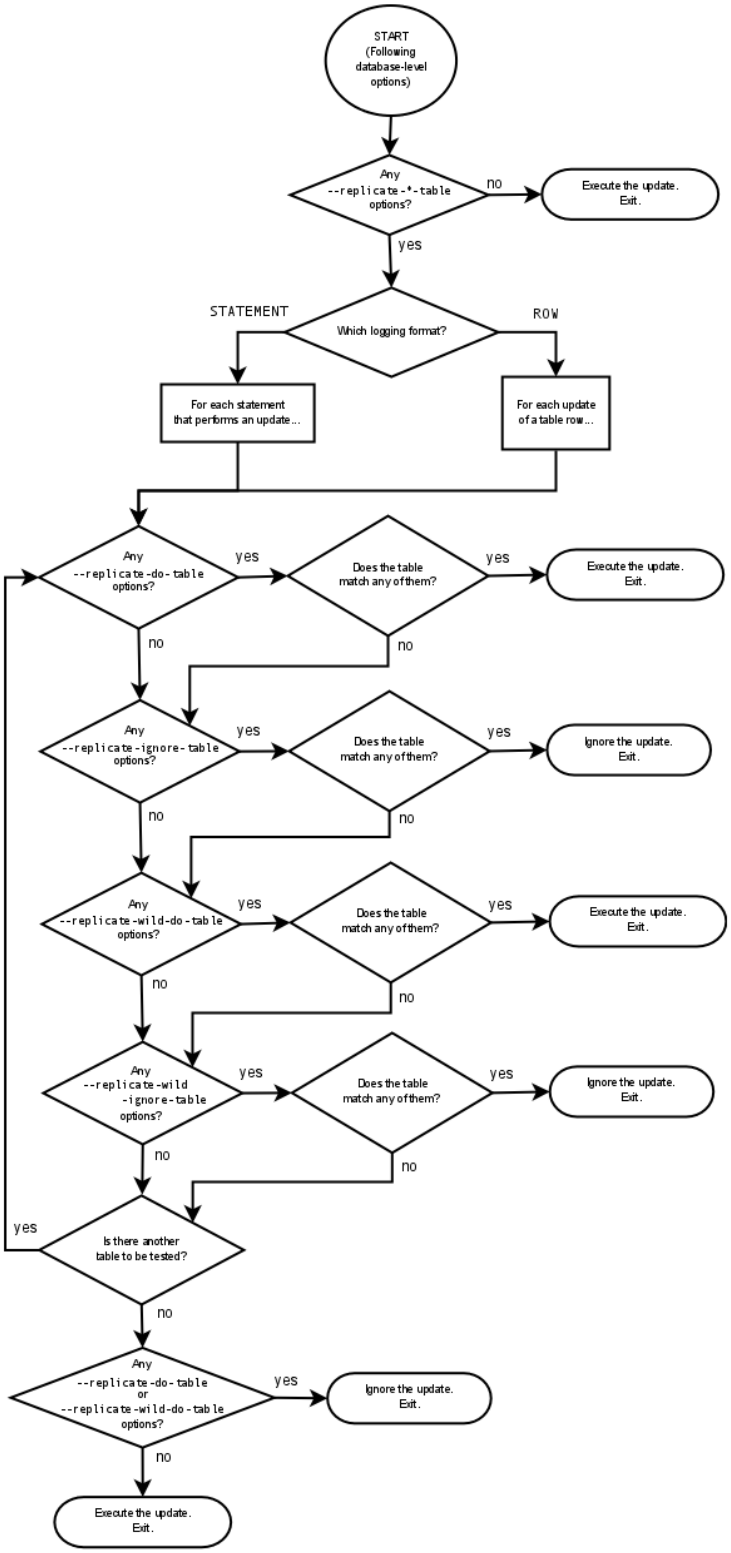
First, as a preliminary condition, the slave checks whether statement-based replication is enabled. If so, and the statement occurs within a stored function, the slave executes the statement and exits.

If row-based replication is enabled, the slave does not know whether a statement occurred within a stored function on the master, so this condition does not apply.

Note

For statement-based replication, replication events represent statements (all changes making up a given event are associated with a single SQL statement); for row-based replication, each event represents a change in a single table row (thus a single statement such as `UPDATE mytable SET mycol = 1` may yield many row-based events). When viewed in terms of events, the process of checking table options is the same for both row-based and statement-based replication.

Having reached this point, if there are no table options, the slave simply executes all events. If there are any `--replicate-do-table` or `--replicate-wild-do-table` options, the event must match one of these if it is to be executed; otherwise, it is ignored. If there are any `--replicate-ignore-table` or `--replicate-wild-ignore-table` options, all events are executed except those that match any of these options. This process is illustrated in the following diagram.



The following steps describe this evaluation in more detail:

1. Are there any table options?

   - **Yes.** Continue to step 2.
   - **No.** Execute the event and exit.

2. Are there any `--replicate-do-table` options?

- **Yes.** Does the table match any of them?
  - **Yes.** Execute the event and exit.
  - **No.** Continue to step 3.
- **No.** Continue to step 3.

3. Are there any `--replicate-ignore-table` options?
   - **Yes.** Does the table match any of them?
     - **Yes.** Ignore the event and exit.
     - **No.** Continue to step 4.
   - **No.** Continue to step 4.
4. Are there any `--replicate-wild-do-table` options?
   - **Yes.** Does the table match any of them?
     - **Yes.** Execute the event and exit.
     - **No.** Continue to step 5.
   - **No.** Continue to step 5.
5. Are there any `--replicate-wild-ignore-table` options?
   - **Yes.** Does the table match any of them?
     - **Yes.** Ignore the event and exit.
     - **No.** Continue to step 6.
   - **No.** Continue to step 6.
6. Are there any `--replicate-do-table` or `--replicate-wild-do-table` options?
   - **Yes.** Ignore the event and exit.
   - **No.** Execute the event and exit.

18.2.5.3 Replication Rule Application

This section provides additional explanation and examples of usage for different combinations of replication filtering options.

Some typical combinations of replication filter rule types are given in the following table:

Condition (Types of Options)	Outcome
No `--replicate-*` options at all:	The slave executes all events that it receives from the master.
`--replicate-*-db` options, but no table options:	The slave accepts or ignores events using the database options. It executes all events permitted by those options because there are no table restrictions.
`--replicate-*-table` options, but no database options:	All events are accepted at the database-checking stage because there are no database conditions. The slave executes or ignores events based solely on the table options.

Condition (Types of Options)	Outcome
A combination of database and table options:	The slave accepts or ignores events using the database options. Then it evaluates all events permitted by those options according to the table options. This can sometimes lead to results that seem counterintuitive, and that may be different depending on whether you are using statement-based or row-based replication; see the text for an example.

A more complex example follows, in which we examine the outcomes for both statement-based and row-based settings.

Suppose that we have two tables `mytbl1` in database `db1` and `mytbl2` in database `db2` on the master, and the slave is running with the following options (and no other replication filtering options):

```
replicate-ignore-db = db1
replicate-do-table  = db2.tbl2
```

Now we execute the following statements on the master:

```
USE db1;
INSERT INTO db2.tbl2 VALUES (1);
```

The results on the slave vary considerably depending on the binary log format, and may not match initial expectations in either case.

Statement-based replication. The `USE` statement causes `db1` to be the default database. Thus the `--replicate-ignore-db` option matches, *and the `INSERT` statement is ignored*. The table options are not checked.

Row-based replication. The default database has no effect on how the slave reads database options when using row-based replication. Thus, the `USE` statement makes no difference in how the `--replicate-ignore-db` option is handled: the database specified by this option does not match the database where the `INSERT` statement changes data, so the slave proceeds to check the table options. The table specified by `--replicate-do-table` matches the table to be updated, *and the row is inserted*.

18.2.5.4 Replication Channel Based Filters

This section explains how to work with replication filters when multiple replication channels exist, for example in a multi-source replication topology. In previous MySQL versions replication filters were global - filters were applied to all replication channels. In MySQL 8.0 replication filters can be global or channel specific, enabling you to configure multi-source replication slaves with replication filters on specific replication channels. For more background information, see Section 18.1.4, “MySQL Multi-Source Replication” and Section 18.2.3, “Replication Channels”.

Important

Group Replication channels cannot be configured with channel specific replication filters.

Overview of Replication Filters and Channels

When multiple replication channels exist, for example in a multi-source replication topology, replication filters are applied as follows:

- Any global replication filter specified is added to the global replication filters of the filter type (`do_db`, `do_ignore_table`, and so on).
- Any channel specific replication filter adds the filter to the specified channel’s replication filters for the specified filter type.

- Each slave replication channel copies global replication filters to its channel specific replication filters if no channel specific replication filter of this type is configured.
- Each channel uses its channel specific replication filters to filter the replication stream.

The syntax to create channel specific replication filters extends the existing SQL statements and command options. When a replication channel is not specified the global replication filter is configured to ensure backwards compatibility. The `CHANGE REPLICATION FILTER` statement supports the `FOR CHANNEL` clause to configure channel specific filters online. The `--replicate-*` command options to configure filters can specify a replication channel using the form `--replicate-filter_type:channel_name=filter_details`. For example, suppose channels `channel_1` and `channel_2` exist before the server starts, starting the slave with the command line options `--replicate-do-db=db1` `--replicate-do-db=channel_1:db2` `--replicate-do-db=db3` `--replicate-ignore-db=db4` `--replicate-ignore-db=channel_2:db5` would result in:

- *Global replication filters*: do_db=db1,db3, ignore_db=db4
- *Channel specific filters on channel_1*: do_db=db2 ignore_db=db4
- *Channel specific filters on channel_2*: do_db=db1,db3 ignore_db=db5

To monitor the replication filters in such a setup use the `replication_applier_global_filters` and `replication_applier_filters` tables.

Configuring Channel Specific Replication Filters at Startup

The replication filter related command options can take an optional *`channel_name`* followed by a colon, followed by the filter specification. The first colon is interpreted as a separator, subsequent colons are interpreted as literal colons. The following command options support channel specific replication filters using this format:

- `--replicate-do-db=`*`channel_name`*`:`*`database_id`*
- `--replicate-ignore-db=`*`channel_name`*`:`*`database_id`*
- `--replicate-do-table=`*`channel_name`*`:`*`table_id`*
- `--replicate-ignore-table=`*`channel_name`*`:`*`table_id`*
- `--replicate-rewrite-db=`*`channel_name`*`:`*`db1-db2`*
- `--replicate-wild-do-table=`*`channel_name`*`:`*`table regexid`*
- `--replicate-wild-ignore-table=`*`channel_name`*`:`*`table regexid`*

If you use a colon but do not specify a *`channel_name`* for the filter option, for example `--replicate-do-db=:`*`database_id`*, the option configures the replication filter for the default replication channel. The default replication channel is the replication channel which always exists once replication has been started, and differs from multi-source replication channels which you create manually. When neither the colon nor a *`channel_name`* is specified the option configures the global replication filters, for example `--replicate-do-db=`*`database_id`* configures the global `--replicate-do-db` filter.

If you configure multiple `rewrite-db=`*`from_name->to_name`* options with the same *`from_name`* database, all filters are added together (put into the `rewrite_do` list) and the first one takes effect.

Changing Channel Specific Replication Filters Online

In addition to the `--replicate-*` options, replication filters can be configured using the `CHANGE REPLICATION FILTER` statement. This removes the need to restart the server, but the slave applier

thread must be stopped while making the change. To make this statement apply the filter to a specific channel, use the `FOR CHANNEL channel` clause. For example:

```
CHANGE REPLICATION FILTER REPLICATE_DO_DB=(db1) FOR CHANNEL channel_1;
```

When a `FOR CHANNEL` clause is provided, the statement acts on the specified channel's replication filters. If multiple types of filters (`do_db`, `do_ignore_table`, `wild_do_table`, and so on) are specified, only the specified filter types are replaced by the statement. In a replication topology with multiple channels, for example on a multi-source replication slave, when no `FOR CHANNEL` clause is provided, the statement acts on the global replication filters and all channels' replication filters, using a similar logic as the `FOR CHANNEL` case. For more information see Section 13.4.2.2, "CHANGE REPLICATION FILTER Syntax".

Removing Channel Specific Replication Filters

When channel specific replication filters have been configured, you can remove the filter by issuing an empty filter type statement. For example to remove all `REPLICATE_REWRITE_DB` filters from a replication channel named `channel_1` issue:

```
CHANGE REPLICATION FILTER REPLICATE_REWRITE_DB=() FOR CHANNEL channel_1;
```

Any `REPLICATE_REWRITE_DB` filters previously configured, using either command options or `CHANGE REPLICATION FILTER`, are removed.

18.3 Replication Solutions

Replication can be used in many different environments for a range of purposes. This section provides general notes and advice on using replication for specific solution types.

For information on using replication in a backup environment, including notes on the setup, backup procedure, and files to back up, see Section 18.3.1, "Using Replication for Backups".

For advice and tips on using different storage engines on the master and slaves, see Section 18.3.4, "Using Replication with Different Master and Slave Storage Engines".

Using replication as a scale-out solution requires some changes in the logic and operation of applications that use the solution. See Section 18.3.5, "Using Replication for Scale-Out".

For performance or data distribution reasons, you may want to replicate different databases to different replication slaves. See Section 18.3.6, "Replicating Different Databases to Different Slaves"

As the number of replication slaves increases, the load on the master can increase and lead to reduced performance (because of the need to replicate the binary log to each slave). For tips on improving your replication performance, including using a single secondary server as a replication master, see Section 18.3.7, "Improving Replication Performance".

For guidance on switching masters, or converting slaves into masters as part of an emergency failover solution, see Section 18.3.8, "Switching Masters During Failover".

To secure your replication communication, you can encrypt the communication channel. For step-by-step instructions, see Section 18.3.9, "Setting Up Replication to Use Secure Connections".

18.3.1 Using Replication for Backups

To use replication as a backup solution, replicate data from the master to a slave, and then back up the data slave. The slave can be paused and shut down without affecting the running operation of the master, so you can produce an effective snapshot of "live" data that would otherwise require the master to be shut down.

How you back up a database depends on its size and whether you are backing up only the data, or the data and the replication slave state so that you can rebuild the slave in the event of failure. There are therefore two choices:

- If you are using replication as a solution to enable you to back up the data on the master, and the size of your database is not too large, the `mysqldump` tool may be suitable. See Section 18.3.1.1, “Backing Up a Slave Using mysqldump”.
- For larger databases, where `mysqldump` would be impractical or inefficient, you can back up the raw data files instead. Using the raw data files option also means that you can back up the binary and relay logs that will enable you to recreate the slave in the event of a slave failure. For more information, see Section 18.3.1.2, “Backing Up Raw Data from a Slave”.

Another backup strategy, which can be used for either master or slave servers, is to put the server in a read-only state. The backup is performed against the read-only server, which then is changed back to its usual read/write operational status. See Section 18.3.1.3, “Backing Up a Master or Slave by Making It Read Only”.

18.3.1.1 Backing Up a Slave Using mysqldump

Using `mysqldump` to create a copy of a database enables you to capture all of the data in the database in a format that enables the information to be imported into another instance of MySQL Server (see Section 4.5.4, “`mysqldump` — A Database Backup Program”). Because the format of the information is SQL statements, the file can easily be distributed and applied to running servers in the event that you need access to the data in an emergency. However, if the size of your data set is very large, `mysqldump` may be impractical.

When using `mysqldump`, you should stop replication on the slave before starting the dump process to ensure that the dump contains a consistent set of data:

1. Stop the slave from processing requests. You can stop replication completely on the slave using `mysqladmin`:

   ```
   shell> mysqladmin stop-slave
   ```

   Alternatively, you can stop only the slave SQL thread to pause event execution:

   ```
   shell> mysql -e 'STOP SLAVE SQL_THREAD;'
   ```

   This enables the slave to continue to receive data change events from the master's binary log and store them in the relay logs using the I/O thread, but prevents the slave from executing these events and changing its data. Within busy replication environments, permitting the I/O thread to run during backup may speed up the catch-up process when you restart the slave SQL thread.

2. Run `mysqldump` to dump your databases. You may either dump all databases or select databases to be dumped. For example, to dump all databases:

   ```
   shell> mysqldump --all-databases > fulldb.dump
   ```

3. Once the dump has completed, start slave operations again:

   ```
   shell> mysqladmin start-slave
   ```

In the preceding example, you may want to add login credentials (user name, password) to the commands, and bundle the process up into a script that you can run automatically each day.

If you use this approach, make sure you monitor the slave replication process to ensure that the time taken to run the backup does not affect the slave's ability to keep up with events from the master. See Section 18.1.7.1, “Checking Replication Status”. If the slave is unable to keep up, you may want to add another slave and distribute the backup process. For an example of how to configure this scenario, see Section 18.3.6, “Replicating Different Databases to Different Slaves”.

18.3.1.2 Backing Up Raw Data from a Slave

To guarantee the integrity of the files that are copied, backing up the raw data files on your MySQL replication slave should take place while your slave server is shut down. If the MySQL server is still

running, background tasks may still be updating the database files, particularly those involving storage engines with background processes such as `InnoDB`. With `InnoDB`, these problems should be resolved during crash recovery, but since the slave server can be shut down during the backup process without affecting the execution of the master it makes sense to take advantage of this capability.

To shut down the server and back up the files:

1. Shut down the slave MySQL server:

   ```
   shell> mysqladmin shutdown
   ```

2. Copy the data files. You can use any suitable copying or archive utility, including `cp`, `tar` or `WinZip`. For example, assuming that the data directory is located under the current directory, you can archive the entire directory as follows:

   ```
   shell> tar cf /tmp/dbbackup.tar ./data
   ```

3. Start the MySQL server again. Under Unix:

   ```
   shell> mysqld_safe &
   ```

   Under Windows:

   ```
   C:\> "C:\Program Files\MySQL\MySQL Server 8.0\bin\mysqld"
   ```

Normally you should back up the entire data directory for the slave MySQL server. If you want to be able to restore the data and operate as a slave (for example, in the event of failure of the slave), then in addition to the slave's data, you should also back up the slave status files, the master info and relay log info repositories, and the relay log files. These files are needed to resume replication after you restore the slave's data.

If you lose the relay logs but still have the `relay-log.info` file, you can check it to determine how far the SQL thread has executed in the master binary logs. Then you can use `CHANGE MASTER TO` with the `MASTER_LOG_FILE` and `MASTER_LOG_POS` options to tell the slave to re-read the binary logs from that point. This requires that the binary logs still exist on the master server.

If your slave is replicating `LOAD DATA INFILE` statements, you should also back up any `SQL_LOAD-*` files that exist in the directory that the slave uses for this purpose. The slave needs these files to resume replication of any interrupted `LOAD DATA INFILE` operations. The location of this directory is the value of the `--slave-load-tmpdir` option. If the server was not started with that option, the directory location is the value of the `tmpdir` system variable.

18.3.1.3 Backing Up a Master or Slave by Making It Read Only

It is possible to back up either master or slave servers in a replication setup by acquiring a global read lock and manipulating the `read_only` system variable to change the read-only state of the server to be backed up:

1. Make the server read-only, so that it processes only retrievals and blocks updates.

2. Perform the backup.

3. Change the server back to its normal read/write state.

Note

The instructions in this section place the server to be backed up in a state that is safe for backup methods that get the data from the server, such as `mysqldump` (see Section 4.5.4, “`mysqldump` — A Database Backup Program”). You should not attempt to use these instructions to make a binary backup by copying files directly because the server may still have modified data cached in memory and not flushed to disk.

The following instructions describe how to do this for a master server and for a slave server. For both scenarios discussed here, suppose that you have the following replication setup:

- A master server M1
- A slave server S1 that has M1 as its master
- A client C1 connected to M1
- A client C2 connected to S1

In either scenario, the statements to acquire the global read lock and manipulate the `read_only` variable are performed on the server to be backed up and do not propagate to any slaves of that server.

Scenario 1: Backup with a Read-Only Master

Put the master M1 in a read-only state by executing these statements on it:

```
mysql> FLUSH TABLES WITH READ LOCK;
mysql> SET GLOBAL read_only = ON;
```

While M1 is in a read-only state, the following properties are true:

- Requests for updates sent by C1 to M1 will block because the server is in read-only mode.
- Requests for query results sent by C1 to M1 will succeed.
- Making a backup on M1 is safe.
- Making a backup on S1 is not safe. This server is still running, and might be processing the binary log or update requests coming from client C2

While M1 is read only, perform the backup. For example, you can use `mysqldump`.

After the backup operation on M1 completes, restore M1 to its normal operational state by executing these statements:

```
mysql> SET GLOBAL read_only = OFF;
mysql> UNLOCK TABLES;
```

Although performing the backup on M1 is safe (as far as the backup is concerned), it is not optimal for performance because clients of M1 are blocked from executing updates.

This strategy applies to backing up a master server in a replication setup, but can also be used for a single server in a nonreplication setting.

Scenario 2: Backup with a Read-Only Slave

Put the slave S1 in a read-only state by executing these statements on it:

```
mysql> FLUSH TABLES WITH READ LOCK;
mysql> SET GLOBAL read_only = ON;
```

While S1 is in a read-only state, the following properties are true:

- The master M1 will continue to operate, so making a backup on the master is not safe.
- The slave S1 is stopped, so making a backup on the slave S1 is safe.

These properties provide the basis for a popular backup scenario: Having one slave busy performing a backup for a while is not a problem because it does not affect the entire network, and the system is still

running during the backup. In particular, clients can still perform updates on the master server, which remains unaffected by backup activity on the slave.

While S1 is read only, perform the backup. For example, you can use `mysqldump`.

After the backup operation on S1 completes, restore S1 to its normal operational state by executing these statements:

```
mysql> SET GLOBAL read_only = OFF;
mysql> UNLOCK TABLES;
```

After the slave is restored to normal operation, it again synchronizes to the master by catching up with any outstanding updates from the binary log of the master.

18.3.2 Handling an Unexpected Halt of a Replication Slave

In order for replication to be resilient to unexpected halts of the server (sometimes described as crash-safe) it must be possible for the slave to recover its state before halting. This section describes the impact of an unexpected halt of a slave during replication and how to configure a slave for the best chance of recovery to continue replication.

After an unexpected halt of a slave, upon restart the I/O thread must recover the information about which transactions have been received, and the SQL thread must recover which transactions have been executed already. For information on the slave logs required for recovery, see Section 18.2.4, “Replication Relay and Status Logs”. The information required for recovery was traditionally stored in files, which had the risk of losing synchrony with the master depending at which stage of processing a transaction the slave halted at, or even corruption of the files themselves. In MySQL 8.0 you can instead use tables to store this information. These tables are created using `InnoDB`, and by using this transactional storage engine the information is always recoverable upon restart. To configure MySQL 8.0 to store the replication information in tables, set `relay_log_info_repository` and `master_info_repository` to `TABLE`. The server then stores information required for the recovery of the I/O thread in the `mysql.slave_master_info` table and information required for the recovery of the SQL thread in the `mysql.slave_relay_log_info` table. With this configuration, DML transactions and also atomic DDLs update the `mysql.slave_relay_log_info` atomically with the replicated data. In all other cases, the `mysql.slave_relay_log_info` table may be missing updates associated with replicated data if the server halts unexpectedly. Restoring updates in this case is a manual process.

Exactly how a replication slave recovers from an unexpected halt is influenced by the chosen method of replication, whether the slave is single-threaded or multi-threaded, the setting of variables such as `relay_log_recovery`, and whether features such as `MASTER_AUTO_POSITION` are being used.

The following table shows the impact of these different factors on how a single-threaded slave recovers from an unexpected halt.

Table 18.5 Factors Influencing Single-threaded Replication Slave Recovery

GTIDs	MASTER_AUTO_POSITION	relay_log_recovery	relay_log_info_repository	Crash type	Recovery guaranteed	Relay log impact
OFF	Any	1	TABLE	Any	Yes	Lost
OFF	Any	1	TABLE	Server	Yes	Lost
OFF	Any	1	Any	OS	No	Lost
OFF	Any	0	TABLE	Server	Yes	Remains
OFF	Any	0	TABLE	OS	No	Remains
ON	ON	Any	Any	Any	Yes	Lost
ON	OFF	0	TABLE	Server	Yes	Remains

GTID	MASTER_AUTO_POSITION	relay_log_recovery	relay_log_info_repository	Crash type	Recovery guaranteed	Relay log impact
ON	OFF	0	Any	OS	No	Remains

As the table shows, when using a single-threaded slave the following configurations are most resilient to unexpected halts:

- When using GTIDs and `MASTER_AUTO_POSITION`, set `relay_log_recovery=0`. With this configuration the setting of `relay_log_info_repository` and other variables does not impact on recovery.
- When using file position based replication, set `relay_log_recovery=1` and `relay_log_info_repository=TABLE`.

Note

During recovery the relay log is lost.

The following table shows the impact of these different factors on how a multi-threaded slave recovers from an unexpected halt.

Table 18.6 Factors Influencing Multi-threaded Replication Slave Recovery

GTID	sync_relay_log	MASTER_AUTO_POSITION	relay_log_recovery	relay_log_info_repository	Crash type	Recovery guaranteed	Relay log impact
OFF	1	Any	1	TABLE	Any	Yes	Lost
OFF	>1	Any	1	TABLE	Server	Yes	Lost
OFF	>1	Any	1	Any	OS	No	Lost
OFF	1	Any	0	TABLE	Server	Yes	Remains
OFF	1	Any	0	TABLE	OS	No	Remains
ON	Any	ON	Any	Any	Any	Yes	Lost
ON	1	OFF	0	TABLE	Server	Yes	Remains
ON	1	OFF	0	Any	OS	No	Remains

As the table shows, when using a multi-threaded slave the following configurations are most resilient to unexpected halts:

- When using GTIDs and `MASTER_AUTO_POSITION`, set `relay_log_recovery=0`. With this configuration the setting of `relay_log_info_repository` and other variables does not impact on recovery.
- When using file position based replication, set `relay_log_recovery=1`, `sync_relay_log=1`, and `relay_log_info_repository=TABLE`.

Note

During recovery the relay log is lost.

It is important to note the impact of `sync_relay_log=1`, which requires a write of to the relay log per transaction. Although this setting is the most resilient to an unexpected halt, with at most one unwritten transaction being lost, it also has the potential to greatly increase the load on storage. Without `sync_relay_log=1`, the effect of an unexpected halt depends on how the relay log is handled by the operating system. Also note that when `relay_log_recovery=0`, the next time the slave is started after an unexpected halt the relay log is processed as part of recovery. After this process completes, the relay log is deleted.

An unexpected halt of a multi-threaded replication slave using the recommended file position based replication configuration above may result in a relay log with transaction inconsistencies (gaps in the sequence of transactions) caused by the unexpected halt. See Section 18.4.1.34, “Replication and Transaction Inconsistencies”. In MySQL 5.7.13 and later, if the relay log recovery process encounters such transaction inconsistencies they are filled and the recovery process continues automatically. In MySQL versions prior to MySQL 5.7.13, this process is not automatic and requires starting the server with `relay_log_recovery=0`, starting the slave with `START SLAVE UNTIL SQL_AFTER_MTS_GAPS` to fix any transaction inconsistencies and then restarting the slave with `relay_log_recovery=1`.

When you are using multi-source replication and `relay_log_recovery=1`, after restarting due to an unexpected halt all replication channels go through the relay log recovery process. Any inconsistencies found in the relay log due to an unexpected halt of a multi-threaded slave are filled.

18.3.3 Monitoring Row-based Replication

The current progress of the replication applier (SQL) thread when using row-based replication is monitored through Performance Schema instrument stages, enabling you to track the processing of operations and check the amount of work completed and work estimated. When these Performance Schema instrument stages are enabled the `events_stages_current` table shows stages for applier threads and their progress. For background information, see Section 25.11.5, “Performance Schema Stage Event Tables”.

To track progress of all three row-based replication event types (write, update, delete):

- Enable the three Performance Schema stages by issuing:

```
mysql> UPDATE performance_schema.setup_instruments SET ENABLED = 'YES'
    -> WHERE NAME LIKE 'stage/sql/Applying batch of row changes%';
```

- Wait for some events to be processed by the replication applier thread and then check progress by looking into the `events_stages_current` table. For example to get progress for `update` events issue:

```
mysql> SELECT WORK_COMPLETED, WORK_ESTIMATED FROM performance_schema.events_stages_current
    -> WHERE EVENT_NAME LIKE 'stage/sql/Applying batch of row changes (update)'
```

- If `binlog_rows_query_log_events` is enabled, information about queries is stored in the binary log and is exposed in the `processlist_info` field. To see the original query that triggered this event:

```
mysql> SELECT db, processlist_state, processlist_info FROM performance_schema.threads
    -> WHERE processlist_state LIKE 'stage/sql/Applying batch of row changes%' AND thread_id = N;
```

18.3.4 Using Replication with Different Master and Slave Storage Engines

It does not matter for the replication process whether the source table on the master and the replicated table on the slave use different engine types. In fact, the `default_storage_engine` system variable is not replicated.

This provides a number of benefits in the replication process in that you can take advantage of different engine types for different replication scenarios. For example, in a typical scale-out scenario (see Section 18.3.5, “Using Replication for Scale-Out”), you want to use `InnoDB` tables on the master to take advantage of the transactional functionality, but use `MyISAM` on the slaves where transaction support is not required because the data is only read. When using replication in a data-logging environment you may want to use the `Archive` storage engine on the slave.

Configuring different engines on the master and slave depends on how you set up the initial replication process:

- If you used `mysqldump` to create the database snapshot on your master, you could edit the dump file text to change the engine type used on each table.

  Another alternative for `mysqldump` is to disable engine types that you do not want to use on the slave before using the dump to build the data on the slave. For example, you can add the `--skip-federated` option on your slave to disable the `FEDERATED` engine. If a specific engine does not exist for a table to be created, MySQL will use the default engine type, usually `MyISAM`. (This requires that the `NO_ENGINE_SUBSTITUTION` SQL mode is not enabled.) If you want to disable additional engines in this way, you may want to consider building a special binary to be used on the slave that only supports the engines you want.

- If you are using raw data files (a binary backup) to set up the slave, you will be unable to change the initial table format. Instead, use `ALTER TABLE` to change the table types after the slave has been started.

- For new master/slave replication setups where there are currently no tables on the master, avoid specifying the engine type when creating new tables.

If you are already running a replication solution and want to convert your existing tables to another engine type, follow these steps:

1. Stop the slave from running replication updates:

   ```
   mysql> STOP SLAVE;
   ```

   This will enable you to change engine types without interruptions.

2. Execute an `ALTER TABLE ... ENGINE='engine_type'` for each table to be changed.

3. Start the slave replication process again:

   ```
   mysql> START SLAVE;
   ```

Although the `default_storage_engine` variable is not replicated, be aware that `CREATE TABLE` and `ALTER TABLE` statements that include the engine specification will be correctly replicated to the slave. For example, if you have a CSV table and you execute:

```
mysql> ALTER TABLE csvtable Engine='MyISAM';
```

The above statement will be replicated to the slave and the engine type on the slave will be converted to `MyISAM`, even if you have previously changed the table type on the slave to an engine other than CSV. If you want to retain engine differences on the master and slave, you should be careful to use the `default_storage_engine` variable on the master when creating a new table. For example, instead of:

```
mysql> CREATE TABLE tablea (columna int) Engine=MyISAM;
```

Use this format:

```
mysql> SET default_storage_engine=MyISAM;
mysql> CREATE TABLE tablea (columna int);
```

When replicated, the `default_storage_engine` variable will be ignored, and the `CREATE TABLE` statement will execute on the slave using the slave's default engine.

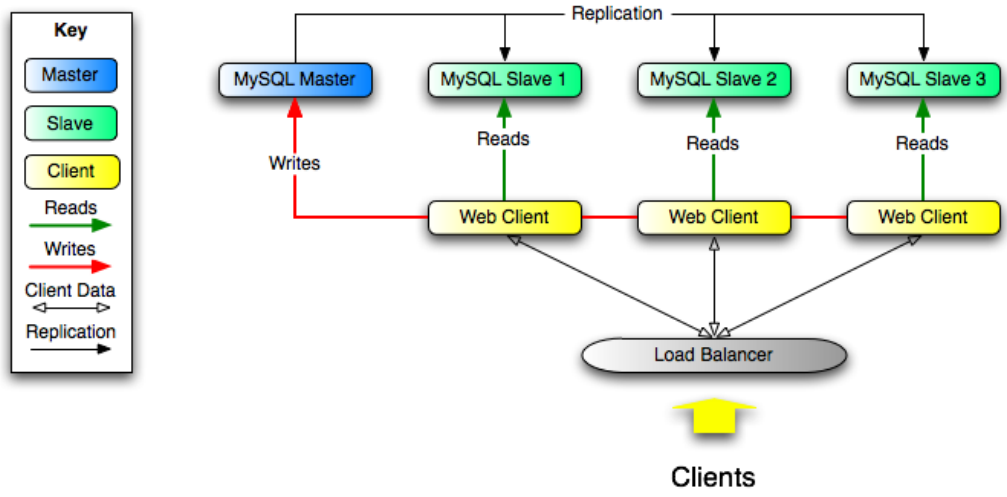
18.3.5 Using Replication for Scale-Out

You can use replication as a scale-out solution; that is, where you want to split up the load of database queries across multiple database servers, within some reasonable limitations.

Because replication works from the distribution of one master to one or more slaves, using replication for scale-out works best in an environment where you have a high number of reads and low number of writes/updates. Most Web sites fit into this category, where users are browsing the Web site, reading articles, posts, or viewing products. Updates only occur during session management, or when making a purchase or adding a comment/message to a forum.

Replication in this situation enables you to distribute the reads over the replication slaves, while still enabling your web servers to communicate with the replication master when a write is required. You can see a sample replication layout for this scenario in Figure 18.1, “Using Replication to Improve Performance During Scale-Out”.

Figure 18.1 Using Replication to Improve Performance During Scale-Out



If the part of your code that is responsible for database access has been properly abstracted/modularized, converting it to run with a replicated setup should be very smooth and easy. Change the implementation of your database access to send all writes to the master, and to send reads to either the master or a slave. If your code does not have this level of abstraction, setting up a replicated system gives you the opportunity and motivation to clean it up. Start by creating a wrapper library or module that implements the following functions:

- `safe_writer_connect()`
- `safe_reader_connect()`
- `safe_reader_statement()`
- `safe_writer_statement()`

`safe_` in each function name means that the function takes care of handling all error conditions. You can use different names for the functions. The important thing is to have a unified interface for connecting for reads, connecting for writes, doing a read, and doing a write.

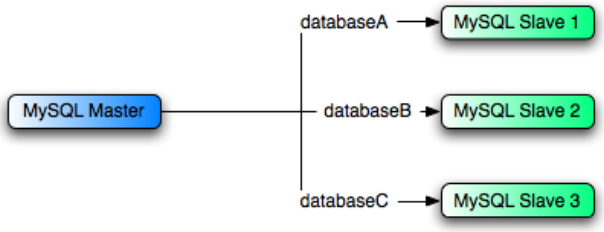
Then convert your client code to use the wrapper library. This may be a painful and scary process at first, but it pays off in the long run. All applications that use the approach just described are able to take advantage of a master/slave configuration, even one involving multiple slaves. The code is much easier to maintain, and adding troubleshooting options is trivial. You need modify only one or two functions; for example, to log how long each statement took, or which statement among those issued gave you an error.

If you have written a lot of code, you may want to automate the conversion task by writing a conversion script. Ideally, your code uses consistent programming style conventions. If not, then you are probably better off rewriting it anyway, or at least going through and manually regularizing it to use a consistent style.

18.3.6 Replicating Different Databases to Different Slaves

There may be situations where you have a single master and want to replicate different databases to different slaves. For example, you may want to distribute different sales data to different departments to help spread the load during data analysis. A sample of this layout is shown in Figure 18.2, “Using Replication to Replicate Databases to Separate Replication Slaves”.

Figure 18.2 Using Replication to Replicate Databases to Separate Replication Slaves



You can achieve this separation by configuring the master and slaves as normal, and then limiting the binary log statements that each slave processes by using the `--replicate-wild-do-table` configuration option on each slave.

Important

You should *not* use `--replicate-do-db` for this purpose when using statement-based replication, since statement-based replication causes this option's affects to vary according to the database that is currently selected. This applies to mixed-format replication as well, since this enables some updates to be replicated using the statement-based format.

However, it should be safe to use `--replicate-do-db` for this purpose if you are using row-based replication only, since in this case the currently selected database has no effect on the option's operation.

For example, to support the separation as shown in Figure 18.2, “Using Replication to Replicate Databases to Separate Replication Slaves”, you should configure each replication slave as follows, before executing `START SLAVE`:

- Replication slave 1 should use `--replicate-wild-do-table=databaseA.%`.
- Replication slave 2 should use `--replicate-wild-do-table=databaseB.%`.
- Replication slave 3 should use `--replicate-wild-do-table=databaseC.%`.

Each slave in this configuration receives the entire binary log from the master, but executes only those events from the binary log that apply to the databases and tables included by the `--replicate-wild-do-table` option in effect on that slave.

If you have data that must be synchronized to the slaves before replication starts, you have a number of choices:

- Synchronize all the data to each slave, and delete the databases, tables, or both that you do not want to keep.
- Use `mysqldump` to create a separate dump file for each database and load the appropriate dump file on each slave.
- Use a raw data file dump and include only the specific files and databases that you need for each slave.

Note

This does not work with `InnoDB` databases unless you use `innodb_file_per_table`.

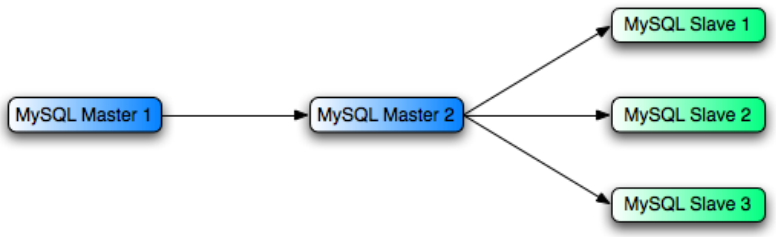
18.3.7 Improving Replication Performance

As the number of slaves connecting to a master increases, the load, although minimal, also increases, as each slave uses a client connection to the master. Also, as each slave must receive a full copy of the master binary log, the network load on the master may also increase and create a bottleneck.

If you are using a large number of slaves connected to one master, and that master is also busy processing requests (for example, as part of a scale-out solution), then you may want to improve the performance of the replication process.

One way to improve the performance of the replication process is to create a deeper replication structure that enables the master to replicate to only one slave, and for the remaining slaves to connect to this primary slave for their individual replication requirements. A sample of this structure is shown in Figure 18.3, “Using an Additional Replication Host to Improve Performance”.

Figure 18.3 Using an Additional Replication Host to Improve Performance



For this to work, you must configure the MySQL instances as follows:

- Master 1 is the primary master where all changes and updates are written to the database. Binary logging should be enabled on this machine.
- Master 2 is the slave to the Master 1 that provides the replication functionality to the remainder of the slaves in the replication structure. Master 2 is the only machine permitted to connect to Master 1. Master 2 also has binary logging enabled, and the `--log-slave-updates` option so that replication instructions from Master 1 are also written to Master 2's binary log so that they can then be replicated to the true slaves.
- Slave 1, Slave 2, and Slave 3 act as slaves to Master 2, and replicate the information from Master 2, which actually consists of the upgrades logged on Master 1.

The above solution reduces the client load and the network interface load on the primary master, which should improve the overall performance of the primary master when used as a direct database solution.

If your slaves are having trouble keeping up with the replication process on the master, there are a number of options available:

- If possible, put the relay logs and the data files on different physical drives. To do this, use the `--relay-log` option to specify the location of the relay log.
- If the slaves are significantly slower than the master, you may want to divide up the responsibility for replicating different databases to different slaves. See Section 18.3.6, “Replicating Different Databases to Different Slaves”.
- If your master makes use of transactions and you are not concerned about transaction support on your slaves, use `MyISAM` or another nontransactional engine on the slaves. See Section 18.3.4, “Using Replication with Different Master and Slave Storage Engines”.

- If your slaves are not acting as masters, and you have a potential solution in place to ensure that you can bring up a master in the event of failure, then you can switch off `--log-slave-updates`. This prevents “dumb” slaves from also logging events they have executed into their own binary log.

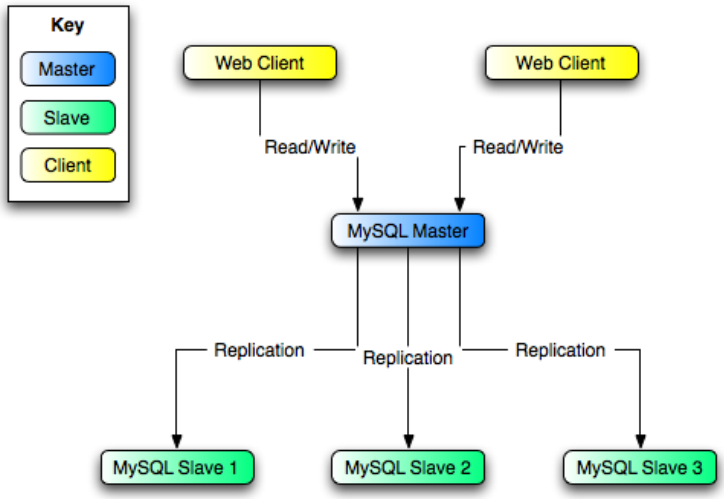
18.3.8 Switching Masters During Failover

When using replication with GTIDs (see Section 18.1.3, “Replication with Global Transaction Identifiers”), you can provide failover between master and slaves in the event of a failure using `mysqlfailover`, which is provided by the MySQL Utilities; see `mysqlfailover` — Automatic replication health monitoring and failover, for more information. If you are not using GTIDs and therefore cannot use `mysqlfailover`, you must set up a master and one or more slaves; then, you need to write an application or script that monitors the master to check whether it is up, and instructs the slaves and applications to change to another master in case of failure. This section discusses some of the issues encountered when setting up failover in this way.

You can tell a slave to change to a new master using the `CHANGE MASTER TO` statement. The slave does not check whether the databases on the master are compatible with those on the slave; it simply begins reading and executing events from the specified coordinates in the new master's binary log. In a failover situation, all the servers in the group are typically executing the same events from the same binary log file, so changing the source of the events should not affect the structure or integrity of the database, provided that you exercise care in making the change.

Slaves should be run with the `--log-bin` option, and if not using GTIDs then they should also be run without `--log-slave-updates`. In this way, the slave is ready to become a master without restarting the slave `mysqld`. Assume that you have the structure shown in Figure 18.4, “Redundancy Using Replication, Initial Structure”.

Figure 18.4 Redundancy Using Replication, Initial Structure



In this diagram, the `MySQL Master` holds the master database, the `MySQL Slave` hosts are replication slaves, and the `Web Client` machines are issuing database reads and writes. Web clients that issue only reads (and would normally be connected to the slaves) are not shown, as they do not need to switch to a new server in the event of failure. For a more detailed example of a read/write scale-out replication structure, see Section 18.3.5, “Using Replication for Scale-Out”.

Each MySQL Slave (`Slave 1`, `Slave 2`, and `Slave 3`) is a slave running with `--log-bin` and without `--log-slave-updates`. Because updates received by a slave from the master are not logged in the binary log unless `--log-slave-updates` is specified, the binary log on each slave is empty initially. If for some reason `MySQL Master` becomes unavailable, you can pick one of the slaves to become the new master. For example, if you pick `Slave 1`, all `Web Clients` should be redirected

to `Slave 1`, which writes the updates to its binary log. `Slave 2` and `Slave 3` should then replicate from `Slave 1`.

The reason for running the slave without `--log-slave-updates` is to prevent slaves from receiving updates twice in case you cause one of the slaves to become the new master. If `Slave 1` has `--log-slave-updates` enabled, it writes any updates that it receives from `Master` in its own binary log. This means that, when `Slave 2` changes from `Master` to `Slave 1` as its master, it may receive updates from `Slave 1` that it has already received from `Master`.

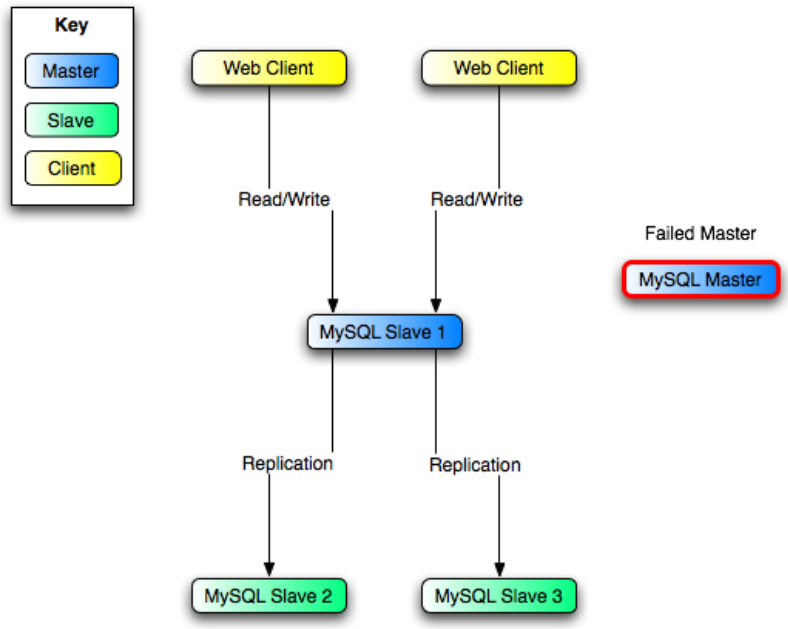
Make sure that all slaves have processed any statements in their relay log. On each slave, issue `STOP SLAVE IO_THREAD`, then check the output of `SHOW PROCESSLIST` until you see `Has read all relay log`. When this is true for all slaves, they can be reconfigured to the new setup. On the slave `Slave 1` being promoted to become the master, issue `STOP SLAVE` and `RESET MASTER`.

On the other slaves `Slave 2` and `Slave 3`, use `STOP SLAVE` and `CHANGE MASTER TO MASTER_HOST='Slave1'` (where `'Slave1'` represents the real host name of `Slave 1`). To use `CHANGE MASTER TO`, add all information about how to connect to `Slave 1` from `Slave 2` or `Slave 3` (*`user`*, *`password`*, *`port`*). When issuing the `CHANGE MASTER TO` statement in this, there is no need to specify the name of the `Slave 1` binary log file or log position to read from, since the first binary log file and position 4, are the defaults. Finally, execute `START SLAVE` on `Slave 2` and `Slave 3`.

Once the new replication setup is in place, you need to tell each `Web Client` to direct its statements to `Slave 1`. From that point on, all updates statements sent by `Web Client` to `Slave 1` are written to the binary log of `Slave 1`, which then contains every update statement sent to `Slave 1` since `Master` died.

The resulting server structure is shown in Figure 18.5, “Redundancy Using Replication, After Master Failure”.

Figure 18.5 Redundancy Using Replication, After Master Failure



When `Master` becomes available again, you should make it a slave of `Slave 1`. To do this, issue on `Master` the same `CHANGE MASTER TO` statement as that issued on `Slave 2` and `Slave 3` previously. `Master` then becomes a slave of `S1ave 1` and picks up the `Web Client` writes that it missed while it was offline.

To make `Master` a master again, use the preceding procedure as if `Slave 1` was unavailable and `Master` was to be the new master. During this procedure, do not forget to run `RESET MASTER` on

`Master` before making `Slave 1`, `Slave 2`, and `Slave 3` slaves of `Master`. If you fail to do this, the slaves may pick up stale writes from the `Web Client` applications dating from before the point at which `Master` became unavailable.

You should be aware that there is no synchronization between slaves, even when they share the same master, and thus some slaves might be considerably ahead of others. This means that in some cases the procedure outlined in the previous example might not work as expected. In practice, however, relay logs on all slaves should be relatively close together.

One way to keep applications informed about the location of the master is to have a dynamic DNS entry for the master. With `bind` you can use `nsupdate` to update the DNS dynamically.

18.3.9 Setting Up Replication to Use Secure Connections

To use a secure connection for encrypting the transfer of the binary log required during replication, both the master and the slave servers must support encrypted network connections. If either server does not support secure connections (because it has not been compiled or configured for them), replication through an encrypted connection is not possible.

Setting up secure connections for replication is similar to doing so for client/server connections. You must obtain (or create) a suitable security certificate that you can use on the master, and a similar certificate (from the same certificate authority) on each slave. You must also obtain suitable key files.

For more information on setting up a server and client for secure connections, see Section 6.4.4, “Configuring MySQL to Use Secure Connections”.

To enable secure connections on the master, you must create or obtain suitable certificate and key files, and then add the following configuration options to the master's configuration within the `[mysqld]` section of the master's `my.cnf` file, changing the file names as necessary:

```
[mysqld]
ssl-ca=cacert.pem
ssl-cert=server-cert.pem
ssl-key=server-key.pem
```

The paths to the files may be relative or absolute; we recommend that you always use complete paths for this purpose.

The options are as follows:

- `ssl-ca` identifies the Certificate Authority (CA) certificate.
- `ssl-cert` identifies the server public key certificate. This can be sent to the client and authenticated against the CA certificate that it has.
- `ssl-key` identifies the server private key.

On the slave, there are two ways to specify the information required for connecting securely to the master. You can either name the slave certificate and key files in the `[client]` section of the slave's `my.cnf` file, or you can explicitly specify that information using the `CHANGE MASTER TO` statement:

- To name the slave certificate and key files using an option file, add the following lines to the `[client]` section of the slave's `my.cnf` file, changing the file names as necessary:

  ```
  [client]
  ssl-ca=cacert.pem
  ssl-cert=client-cert.pem
  ssl-key=client-key.pem
  ```

  Restart the slave server, using the `--skip-slave-start` option to prevent the slave from connecting to the master. Use `CHANGE MASTER TO` to specify the master configuration, using the `MASTER_SSL` option to connect securely:

```
mysql> CHANGE MASTER TO
    -> MASTER_HOST='master_hostname',
    -> MASTER_USER='replicate',
    -> MASTER_PASSWORD='password',
    -> MASTER_SSL=1;
```

- To specify the certificate and key names using the `CHANGE MASTER TO` statement, append the appropriate `MASTER_SSL_xxx` options:

```
mysql> CHANGE MASTER TO
    -> MASTER_HOST='master_hostname',
    -> MASTER_USER='replicate',
    -> MASTER_PASSWORD='password',
    -> MASTER_SSL=1,
    -> MASTER_SSL_CA = 'ca_file_name',
    -> MASTER_SSL_CAPATH = 'ca_directory_name',
    -> MASTER_SSL_CERT = 'cert_file_name',
    -> MASTER_SSL_KEY = 'key_file_name';
```

After the master information has been updated, start the slave replication process:

```
mysql> START SLAVE;
```

You can use the `SHOW SLAVE STATUS` statement to confirm that a secure connection was established successfully.

For more information on the `CHANGE MASTER TO` statement, see Section 13.4.2.1, “CHANGE MASTER TO Syntax”.

If you want to enforce the use of secure connections during replication, create a user and use the `REQUIRE SSL` option, then grant that user the `REPLICATION SLAVE` privilege. For example:

```
mysql> CREATE USER 'repl'@'%.mydomain.com' IDENTIFIED BY 'slavepass'
    -> REQUIRE SSL;
mysql> GRANT REPLICATION SLAVE ON *.*
    -> TO 'repl'@'%.mydomain.com';
```

If the account already exists, you can add `REQUIRE SSL` to it with this statement:

```
mysql> ALTER USER 'repl'@'%.mydomain.com' REQUIRE SSL;
```

18.3.10 Semisynchronous Replication

In addition to the built-in asynchronous replication, MySQL 8.0 supports an interface to semisynchronous replication that is implemented by plugins. This section discusses what semisynchronous replication is and how it works. The following sections cover the administrative interface to semisynchronous replication and how to install, configure, and monitor it.

MySQL replication by default is asynchronous. The master writes events to its binary log but does not know whether or when a slave has retrieved and processed them. With asynchronous replication, if the master crashes, transactions that it has committed might not have been transmitted to any slave. Consequently, failover from master to slave in this case may result in failover to a server that is missing transactions relative to the master.

Semisynchronous replication can be used as an alternative to asynchronous replication:

- A slave indicates whether it is semisynchronous-capable when it connects to the master.
- If semisynchronous replication is enabled on the master side and there is at least one semisynchronous slave, a thread that performs a transaction commit on the master blocks and

waits until at least one semisynchronous slave acknowledges that it has received all events for the transaction, or until a timeout occurs.

- The slave acknowledges receipt of a transaction's events only after the events have been written to its relay log and flushed to disk.
- If a timeout occurs without any slave having acknowledged the transaction, the master reverts to asynchronous replication. When at least one semisynchronous slave catches up, the master returns to semisynchronous replication.
- Semisynchronous replication must be enabled on both the master and slave sides. If semisynchronous replication is disabled on the master, or enabled on the master but on no slaves, the master uses asynchronous replication.

While the master is blocking (waiting for acknowledgment from a slave), it does not return to the session that performed the transaction. When the block ends, the master returns to the session, which then can proceed to execute other statements. At this point, the transaction has committed on the master side, and receipt of its events has been acknowledged by at least one slave.

The number of slave acknowledgments the master must receive per transaction before proceeding is configurable using the `rpl_semi_sync_master_wait_for_slave_count` system variable. The default value is 1.

Blocking also occurs after rollbacks that are written to the binary log, which occurs when a transaction that modifies nontransactional tables is rolled back. The rolled-back transaction is logged even though it has no effect for transactional tables because the modifications to the nontransactional tables cannot be rolled back and must be sent to slaves.

For statements that do not occur in transactional context (that is, when no transaction has been started with `START TRANSACTION` or `SET autocommit = 0`), autocommit is enabled and each statement commits implicitly. With semisynchronous replication, the master blocks for each such statement, just as it does for explicit transaction commits.

To understand what the “semi” in “semisynchronous replication” means, compare it with asynchronous and fully synchronous replication:

- With asynchronous replication, the master writes events to its binary log and slaves request them when they are ready. There is no guarantee that any event will ever reach any slave.
- With fully synchronous replication, when a master commits a transaction, all slaves also will have committed the transaction before the master returns to the session that performed the transaction. The drawback of this is that there might be a lot of delay to complete a transaction.
- Semisynchronous replication falls between asynchronous and fully synchronous replication. The master waits only until at least one slave has received and logged the events. It does not wait for all slaves to acknowledge receipt, and it requires only receipt, not that the events have been fully executed and committed on the slave side.

Compared to asynchronous replication, semisynchronous replication provides improved data integrity because when a commit returns successfully, it is known that the data exists in at least two places. Until a semisynchronous master receives acknowledgment from the number of slaves configured by `rpl_semi_sync_master_wait_for_slave_count`, the transaction is on hold and not committed.

Semisynchronous replication also places a rate limit on busy sessions by constraining the speed at which binary log events can be sent from master to slave. When one user is too busy, this will slow it down, which is useful in some deployment situations.

Semisynchronous replication does have some performance impact because commits are slower due to the need to wait for slaves. This is the tradeoff for increased data integrity. The amount of slowdown is at least the TCP/IP roundtrip time to send the commit to the slave and wait for the acknowledgment of receipt by the slave. This means that semisynchronous replication works best for close servers communicating over fast networks, and worst for distant servers communicating over slow networks.

The `rpl_semi_sync_master_wait_point` system variable controls the point at which a semisynchronous replication master waits for slave acknowledgment of transaction receipt before returning a status to the client that committed the transaction. These values are permitted:

- `AFTER_SYNC` (the default): The master writes each transaction to its binary log and the slave, and syncs the binary log to disk. The master waits for slave acknowledgment of transaction receipt after the sync. Upon receiving acknowledgment, the master commits the transaction to the storage engine and returns a result to the client, which then can proceed.
- `AFTER_COMMIT`: The master writes each transaction to its binary log and the slave, syncs the binary log, and commits the transaction to the storage engine. The master waits for slave acknowledgment of transaction receipt after the commit. Upon receiving acknowledgment, the master returns a result to the client, which then can proceed.

The replication characteristics of these settings differ as follows:

- With `AFTER_SYNC`, all clients see the committed transaction at the same time: After it has been acknowledged by the slave and committed to the storage engine on the master. Thus, all clients see the same data on the master.

  In the event of master failure, all transactions committed on the master have been replicated to the slave (saved to its relay log). A crash of the master and failover to the slave is lossless because the slave is up to date.

- With `AFTER_COMMIT`, the client issuing the transaction gets a return status only after the server commits to the storage engine and receives slave acknowledgment. After the commit and before slave acknowledgment, other clients can see the committed transaction before the committing client.

  If something goes wrong such that the slave does not process the transaction, then in the event of a master crash and failover to the slave, it is possible that such clients will see a loss of data relative to what they saw on the master.

18.3.10.1 Semisynchronous Replication Administrative Interface

The administrative interface to semisynchronous replication has several components:

- Two plugins implement semisynchronous capability. There is one plugin for the master side and one for the slave side.
- System variables control plugin behavior. Some examples:
  - `rpl_semi_sync_master_enabled`

    Controls whether semisynchronous replication is enabled on the master. To enable or disable the plugin, set this variable to 1 or 0, respectively. The default is 0 (off).

  - `rpl_semi_sync_master_timeout`

    A value in milliseconds that controls how long the master waits on a commit for acknowledgment from a slave before timing out and reverting to asynchronous replication. The default value is 10000 (10 seconds).

  - `rpl_semi_sync_slave_enabled`

    Similar to `rpl_semi_sync_master_enabled`, but controls the slave plugin.

  All `rpl_semi_sync_`*`xxx`* system variables are described at Section 5.1.5, “Server System Variables”.

- Status variables enable semisynchronous replication monitoring. Some examples:
  - `Rpl_semi_sync_master_clients`

The number of semisynchronous slaves.

- `Rpl_semi_sync_master_status`

  Whether semisynchronous replication currently is operational on the master. The value is 1 if the plugin has been enabled and a commit acknowledgment has not occurred. It is 0 if the plugin is not enabled or the master has fallen back to asynchronous replication due to commit acknowledgment timeout.

- `Rpl_semi_sync_master_no_tx`

  The number of commits that were not acknowledged successfully by a slave.

- `Rpl_semi_sync_master_yes_tx`

  The number of commits that were acknowledged successfully by a slave.

- `Rpl_semi_sync_slave_status`

  Whether semisynchronous replication currently is operational on the slave. This is 1 if the plugin has been enabled and the slave I/O thread is running, 0 otherwise.

All `Rpl_semi_sync_xxx` status variables are described at Section 5.1.7, “Server Status Variables”.

The system and status variables are available only if the appropriate master or slave plugin has been installed with `INSTALL PLUGIN`.

18.3.10.2 Semisynchronous Replication Installation and Configuration

Semisynchronous replication is implemented using plugins, so the plugins must be installed into the server to make them available. After a plugin has been installed, you control it by means of the system variables associated with it. These system variables are unavailable until the associated plugin has been installed.

This section describes how to install the semisynchronous replication plugins. For general information about installing plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To use semisynchronous replication, the following requirements must be satisfied:

- MySQL 5.5 or higher must be installed.
- The capability of installing plugins requires a MySQL server that supports dynamic loading. To verify this, check that the value of the `have_dynamic_loading` system variable is `YES`. Binary distributions should support dynamic loading.
- Replication must already be working, see Section 18.1, “Configuring Replication”.
- There must not be multiple replication channels configured. Semisynchronous replication is only compatible with the default replication channel. See Section 18.2.3, “Replication Channels”.

To set up semisynchronous replication, use the following instructions. The `INSTALL PLUGIN`, `SET GLOBAL`, `STOP SLAVE`, and `START SLAVE` statements mentioned here require the `REPLICATION_SLAVE_ADMIN` or `SUPER` privilege.

MySQL distributions include semisynchronous replication plugin files for the master side and the slave side.

To be usable by a master or slave server, the appropriate plugin library file must be located in the MySQL plugin directory (the directory named by the `plugin_dir` system variable). If necessary, set the value of `plugin_dir` at server startup to tell the server the plugin directory location.

The plugin library file base names are `semisync_master` and `semisync_slave`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

The master plugin library file must be present in the plugin directory of the master server. The slave plugin library file must be present in the plugin directory of each slave server.

To load the plugins, use the `INSTALL PLUGIN` statement on the master and on each slave that is to be semisynchronous (adjust the `.so` suffix for your platform as necessary).

On the master:

```
INSTALL PLUGIN rpl_semi_sync_master SONAME 'semisync_master.so';
```

On each slave:

```
INSTALL PLUGIN rpl_semi_sync_slave SONAME 'semisync_slave.so';
```

If an attempt to install a plugin results in an error on Linux similar to that shown here, you must install `libimf`:

```
mysql> INSTALL PLUGIN rpl_semi_sync_master SONAME 'semisync_master.so';
ERROR 1126 (HY000): Can't open shared library
'/usr/local/mysql/lib/plugin/semisync_master.so'
(errno: 22 libimf.so: cannot open shared object file:
No such file or directory)
```

You can obtain `libimf` from http://dev.mysql.com/downloads/os-linux.html.

To see which plugins are installed, use the `SHOW PLUGINS` statement, or query the `INFORMATION_SCHEMA.PLUGINS` table.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE '%semi%';
+----------------------+---------------+
| PLUGIN_NAME          | PLUGIN_STATUS |
+----------------------+---------------+
| rpl_semi_sync_master | ACTIVE        |
+----------------------+---------------+
```

If the plugin fails to initialize, check the server error log for diagnostic messages.

After a semisynchronous replication plugin has been installed, it is disabled by default. The plugins must be enabled both on the master side and the slave side to enable semisynchronous replication. If only one side is enabled, replication will be asynchronous.

To control whether an installed plugin is enabled, set the appropriate system variables. You can set these variables at runtime using `SET GLOBAL`, or at server startup on the command line or in an option file.

At runtime, these master-side system variables are available:

```
SET GLOBAL rpl_semi_sync_master_enabled = {0|1};
SET GLOBAL rpl_semi_sync_master_timeout = N;
```

On the slave side, this system variable is available:

```
SET GLOBAL rpl_semi_sync_slave_enabled = {0|1};
```

For `rpl_semi_sync_master_enabled` or `rpl_semi_sync_slave_enabled`, the value should be 1 to enable semisynchronous replication or 0 to disable it. By default, these variables are set to 0.

For `rpl_semi_sync_master_timeout`, the value *`N`* is given in milliseconds. The default value is 10000 (10 seconds).

If you enable semisynchronous replication on a slave at runtime, you must also start the slave I/O thread (stopping it first if it is already running) to cause the slave to connect to the master and register as a semisynchronous slave:

```
STOP SLAVE IO_THREAD;
START SLAVE IO_THREAD;
```

If the I/O thread is already running and you do not restart it, the slave continues to use asynchronous replication.

At server startup, the variables that control semisynchronous replication can be set as command-line options or in an option file. A setting listed in an option file takes effect each time the server starts. For example, you can set the variables in `my.cnf` files on the master and slave sides as follows.

On the master:

```
[mysqld]
rpl_semi_sync_master_enabled=1
rpl_semi_sync_master_timeout=1000 # 1 second
```

On each slave:

```
[mysqld]
rpl_semi_sync_slave_enabled=1
```

18.3.10.3 Semisynchronous Replication Monitoring

The plugins for the semisynchronous replication capability expose several system and status variables that you can examine to determine its configuration and operational state.

The system variable reflect how semisynchronous replication is configured. To check their values, use `SHOW VARIABLES`:

```
mysql> SHOW VARIABLES LIKE 'rpl_semi_sync%';
```

The status variables enable you to monitor the operation of semisynchronous replication. To check their values, use `SHOW STATUS`:

```
mysql> SHOW STATUS LIKE 'Rpl_semi_sync%';
```

When the master switches between asynchronous or semisynchronous replication due to commit-blocking timeout or a slave catching up, it sets the value of the `Rpl_semi_sync_master_status` status variable appropriately. Automatic fallback from semisynchronous to asynchronous replication on the master means that it is possible for the `rpl_semi_sync_master_enabled` system variable to have a value of 1 on the master side even when semisynchronous replication is in fact not operational at the moment. You can monitor the `Rpl_semi_sync_master_status` status variable to determine whether the master currently is using asynchronous or semisynchronous replication.

To see how many semisynchronous slaves are connected, check `Rpl_semi_sync_master_clients`.

The number of commits that have been acknowledged successfully or unsuccessfully by slaves are indicated by the `Rpl_semi_sync_master_yes_tx` and `Rpl_semi_sync_master_no_tx` variables.

On the slave side, `Rpl_semi_sync_slave_status` indicates whether semisynchronous replication currently is operational.

18.3.11 Delayed Replication

MySQL supports delayed replication such that a slave server deliberately executes transactions later than the master by at least a specified amount of time. This section describes how to configure a replication delay on a slave, and how to monitor replication delay.

Compared to previous versions, MySQL 8.0 uses a different method for delaying replication that depends on two timestamps added in MySQL 8.0.1 - `immediate_commit_timestamp` and `original_commit_timestamp`, which are explained in Replication Delay Timestamps . If all servers in the replication topology are running MySQL version 8.0.1 or later, delayed replication is measured using these timestamps. If either the immediate master or slave is not using these timestamps, the implementation of delayed replication from MySQL 5.7 used, see Delayed Replication. This section describes delayed replication between servers which are all using these timestamps.

The default replication delay is 0 seconds. Use the `CHANGE MASTER TO MASTER_DELAY=N` statement to set the delay to *`N`* seconds. A transaction received from the master is not executed until at least *`N`* seconds later than its commit on the immediate master. The delay happens per transaction (not event as in previous MySQL versions) and the actual delay is imposed only on `gtid_log_event` or `anonymous_gtid_log_event`. The other events in the transaction always follow these events without any waiting time imposed on them.

Note

`START SLAVE` and `STOP SLAVE` take effect immediately and ignore any delay. `RESET SLAVE` resets the delay to 0.

The `replication_applier_configuration` Performance Schema table contains the `DESIRED_DELAY` column which shows the delay configured using the `MASTER_DELAY` option. The `replication_applier_status` Performance Schema table contains the `REMAINING_DELAY` column which shows the number of delay seconds remaining.

Delayed replication can be used for several purposes:

- To protect against user mistakes on the master. With a delay you can roll back a delayed slave to the time just before the mistake.
- To test how the system behaves when there is a lag. For example, in an application, a lag might be caused by a heavy load on the slave. However, it can be difficult to generate this load level. Delayed replication can simulate the lag without having to simulate the load. It can also be used to debug conditions related to a lagging slave.
- To inspect what the database looked like in the past, without having to reload a backup. For example, by configuring a slave with a delay of one week, if you then need to see what the database looked like before the last few days' worth of development, the delayed slave can be inspected.

Replication Delay Timestamps

MySQL 8.0 provides a new method for measuring delay (also referred to as replication lag) in replication topologies that depends on the following timestamps associated with the GTID of each transaction (instead of each event) written to the binary log.

- `original_commit_timestamp`: the number of microseconds since epoch when the transaction was written (committed) to the binary log of the original master.

- `immediate_commit_timestamp`: the number of microseconds since epoch when the transaction was written (committed) to the binary log of the immediate master.

The output of `mysqlbinlog` displays these timestamps in two formats, microseconds from epoch and also `TIMESTAMP` format, which is based on the user defined timezone for better readability. For example:

```
#170404 10:48:05 server id 1  end_log_pos 233 CRC32 0x016ce647     GTID    last_committed=0
\ sequence_number=1    original_committed_timestamp=1491299285661130    immediate_commit_timestamp=1491
# original_commit_timestamp=1491299285661130 (2017-04-04 10:48:05.661130 WEST)
# immediate_commit_timestamp=1491299285843771 (2017-04-04 10:48:05.843771 WEST)
 /*!80001 SET @@session.original_commit_timestamp=1491299285661130*//*!*/;
   SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa-aaaa-aaaa-aaaaaaaaaaaa:1'/*!*/;
# at 233
```

As a rule, the `original_commit_timestamp` is always the same on all replicas where the transaction is applied. In master-slave replication, the `original_commit_timestamp` of a transaction in the (original) master's binary log is always the same as its `immediate_commit_timestamp`. In the slave's relay log, the `original_commit_timestamp` and `immediate_commit_timestamp` of the transaction are the same as in the master's binary log; whereas in its own binary log, the transaction's `immediate_commit_timestamp` corresponds to when the slave committed the transaction.

In a Group Replication setup, when the original master is a member of a group, the `original_commit_timestamp` is generated when the transaction is ready to be committed. In other words, when it finished executing on the original master and its write set is ready to be sent to all members of the group for certification. Therefore, the same `original_commit_timestamp` is replicated to all servers (regardless of whether it is a group member or slave replicating from a member) applying the transaction and each stores in its binary log the local commit time using `immediate_commit_timestamp`.

View change events, which are exclusive to Group Replication, are a special case. Transactions containing these events are generated by each server but share the same GTID (so, they are not first executed in a master and then replicated to the group, but all members of the group execute and apply the same transaction). Since there is no original master, these transactions have their `original_commit_timestamp` set to zero.

Monitoring Replication Delay

One of the most common ways to monitor replication delay (lag) in previous MySQL versions was by relying on the `Seconds_Behind_Master` field in the output of `SHOW SLAVE STATUS`. However, this metric is not suitable when using replication topologies more complex than the traditional master-slave setup, such as Group Replication. The addition of `immediate_commit_timestamp` and `original_commit_timestamp` to MySQL 8 provides a much finer degree of information about replication delay. The recommended method to monitor replication delay in a topology that supports these timestamps is using the following Performance Schema tables.

- `replication_connection_status`: current status of the connection to the master, provides information on the last and current transaction the connection thread queued into the relay log.
- `replication_applier_status_by_coordinator`: current status of the coordinator thread that only displays information when using a multi-threaded slave, provides information on the last transaction buffered by the coordinator thread to a worker's queue, as well as the transaction it is currently buffering.
- `replication_applier_status_by_worker`: current status of the thread(s) applying transactions received from the master, provides information about the transactions applied by the applier thread, or by each worker when using a multi-threaded slave.

Using these tables you can monitor information about the last transaction the corresponding thread processed and the transaction that thread is currently processing. This information comprises:

- a transaction’s GTID
- a transaction's `original_commit_timestamp` and `immediate_commit_timestamp`, retrieved from the slave’s relay log
- the time a thread started processing a transaction
- for the last processed transaction, the time the thread finished processing it

In addition to the Performance Schema tables, the output of `SHOW SLAVE STATUS` has three fields that show:

- `SQL_Delay`: A nonnegative integer indicating the replication delay configured using `CHANGE MASTER TO MASTER_DELAY=N`, measured in seconds.
- `SQL_Remaining_Delay`: When `Slave_SQL_Running_State` is `Waiting until MASTER_DELAY seconds after master executed event`, this field contains an integer indicating the number of seconds left of the delay. At other times, this field is `NULL`.
- `Slave_SQL_Running_State`: A string indicating the state of the SQL thread (analogous to `Slave_IO_State`). The value is identical to the `State` value of the SQL thread as displayed by `SHOW PROCESSLIST`.

When the slave SQL thread is waiting for the delay to elapse before executing an event, `SHOW PROCESSLIST` displays its `State` value as `Waiting until MASTER_DELAY seconds after master executed event`.

18.4 Replication Notes and Tips

18.4.1 Replication Features and Issues

The following sections provide information about what is supported and what is not in MySQL replication, and about specific issues and situations that may occur when replicating certain statements.

Statement-based replication depends on compatibility at the SQL level between the master and slave. In others, successful SBR requires that any SQL features used be supported by both the master and the slave servers. For example, if you use a feature on the master server that is available only in MySQL 8.0 (or later), you cannot replicate to a slave that uses MySQL 5.7 (or earlier).

Such incompatibilities also can occur within a release series when using pre-production releases of MySQL. For example, the `SLEEP()` function is available beginning with MySQL 5.0.12. If you use this function on the master, you cannot replicate to a slave that uses MySQL 5.0.11 or earlier.

For this reason, use Generally Available (GA) releases of MySQL for statement-based replication in a production setting, since we do not introduce new SQL statements or change their behavior within a given release series once that series reaches GA release status.

If you are planning to use statement-based replication between MySQL 8.0 and a previous MySQL release series, it is also a good idea to consult the edition of the *MySQL Reference Manual* corresponding to the earlier release series for information regarding the replication characteristics of that series.

With MySQL's statement-based replication, there may be issues with replicating stored routines or triggers. You can avoid these issues by using MySQL's row-based replication instead. For a detailed list of issues, see Section 23.7, “Binary Logging of Stored Programs”. For more information about row-based logging and row-based replication, see Section 5.4.4.1, “Binary Logging Formats”, and Section 18.2.1, “Replication Formats”.

For additional information specific to replication and `InnoDB`, see Section 15.18, “InnoDB and MySQL Replication”. For information relating to replication with MySQL Cluster, see NDB Cluster Replication.

18.4.1.1 Replication and AUTO_INCREMENT

Statement-based replication of `AUTO_INCREMENT`, `LAST_INSERT_ID()`, and `TIMESTAMP` values is done correctly, subject to the following exceptions:

- A statement invoking a trigger or function that causes an update to an `AUTO_INCREMENT` column is not replicated correctly using statement-based replication. In MySQL 8.0, such statements are marked as unsafe. (Bug #45677)
- An `INSERT` into a table that has a composite primary key that includes an `AUTO_INCREMENT` column that is not the first column of this composite key is not safe for statement-based logging or replication. In MySQL 8.0 and later, such statements are marked as unsafe. (Bug #11754117, Bug #45670)

  This issue does not affect tables using the `InnoDB` storage engine, since an `InnoDB` table with an AUTO_INCREMENT column requires at least one key where the auto-increment column is the only or leftmost column.
- Adding an `AUTO_INCREMENT` column to a table with `ALTER TABLE` might not produce the same ordering of the rows on the slave and the master. This occurs because the order in which the rows are numbered depends on the specific storage engine used for the table and the order in which the rows were inserted. If it is important to have the same order on the master and slave, the rows must be ordered before assigning an `AUTO_INCREMENT` number. Assuming that you want to add an `AUTO_INCREMENT` column to a table `t1` that has columns `col1` and `col2`, the following statements produce a new table `t2` identical to `t1` but with an `AUTO_INCREMENT` column:

```
CREATE TABLE t2 LIKE t1;
ALTER TABLE t2 ADD id INT AUTO_INCREMENT PRIMARY KEY;
INSERT INTO t2 SELECT * FROM t1 ORDER BY col1, col2;
```

Important

To guarantee the same ordering on both master and slave, the `ORDER BY` clause must name *all* columns of `t1`.

The instructions just given are subject to the limitations of `CREATE TABLE ... LIKE`: Foreign key definitions are ignored, as are the `DATA DIRECTORY` and `INDEX DIRECTORY` table options. If a table definition includes any of those characteristics, create `t2` using a `CREATE TABLE` statement that is identical to the one used to create `t1`, but with the addition of the `AUTO_INCREMENT` column.

Regardless of the method used to create and populate the copy having the `AUTO_INCREMENT` column, the final step is to drop the original table and then rename the copy:

```
DROP t1;
ALTER TABLE t2 RENAME t1;
```

See also Section B.5.6.1, “Problems with ALTER TABLE”.

18.4.1.2 Replication and BLACKHOLE Tables

The `BLACKHOLE` storage engine accepts data but discards it and does not store it. When performing binary logging, all inserts to such tables are always logged, regardless of the logging format in use. Updates and deletes are handled differently depending on whether statement based or row based logging is in use. With the statement based logging format, all statements affecting `BLACKHOLE` tables are logged, but their effects ignored. When using row-based logging, updates and deletes to such tables are simply skipped—they are not written to the binary log. A warning is logged whenever this occurs.

For this reason we recommend when you replicate to tables using the `BLACKHOLE` storage engine that you have the `binlog_format` server variable set to `STATEMENT`, and not to either `ROW` or `MIXED`.

18.4.1.3 Replication and Character Sets

The following applies to replication between MySQL servers that use different character sets:

- If the master has databases with a character set different from the global `character_set_server` value, you should design your `CREATE TABLE` statements so that they do not implicitly rely on the database default character set. A good workaround is to state the character set and collation explicitly in `CREATE TABLE` statements.

18.4.1.4 Replication and CHECKSUM TABLE

`CHECKSUM TABLE` returns a checksum that is calculated row by row, using a method that depends on the table row storage format, which is not guaranteed to remain the same between MySQL release series. For example, the storage format for temporal types such as `TIME`, `DATETIME`, and `TIMESTAMP` changed in MySQL 5.6 prior to MySQL 5.6.5, so if a 5.5 table is upgraded to MySQL 5.6, the checksum value may change.

18.4.1.5 Replication of CREATE ... IF NOT EXISTS Statements

MySQL applies these rules when various `CREATE ... IF NOT EXISTS` statements are replicated:

- Every `CREATE DATABASE IF NOT EXISTS` statement is replicated, whether or not the database already exists on the master.
- Similarly, every `CREATE TABLE IF NOT EXISTS` statement without a `SELECT` is replicated, whether or not the table already exists on the master. This includes `CREATE TABLE IF NOT EXISTS ... LIKE`. Replication of `CREATE TABLE IF NOT EXISTS ... SELECT` follows somewhat different rules; see Section 18.4.1.6, “Replication of CREATE TABLE ... SELECT Statements”, for more information.
- `CREATE EVENT IF NOT EXISTS` is always replicated in MySQL 8.0, whether or not the event named in the statement already exists on the master.

See also Bug #45574.

18.4.1.6 Replication of CREATE TABLE ... SELECT Statements

This section discusses how MySQL replicates `CREATE TABLE ... SELECT` statements.

MySQL 8.0 does not allow a `CREATE TABLE ... SELECT` statement to make any changes in tables other than the table that is created by the statement. Some older versions of MySQL permitted these statements to do so; this means that, when using statement-based replication between a MySQL 5.6 or later slave and a master running a previous version of MySQL, a `CREATE TABLE ... SELECT` statement causing changes in other tables on the master fails on the slave, causing replication to stop. To prevent this from happening, you should use row-based replication, rewrite the offending statement before running it on the master, or upgrade the master to MySQL 8.0. (If you choose to upgrade the master, keep in mind that such a `CREATE TABLE ... SELECT` statement fails following the upgrade unless it is rewritten to remove any side effects on other tables.) This is not an issue when using row-based replication, because the statement is logged as a `CREATE TABLE` statement with any changes to table data logged as row-insert events, rather than as the entire `CREATE TABLE ... SELECT`.

These behaviors are not dependent on MySQL version:

- `CREATE TABLE ... SELECT` always performs an implicit commit (Section 13.3.3, “Statements That Cause an Implicit Commit”).
- If destination table does not exist, logging occurs as follows. It does not matter whether `IF NOT EXISTS` is present.
  - `STATEMENT` or `MIXED` format: The statement is logged as written.

  - `ROW` format: The statement is logged as a `CREATE TABLE` statement followed by a series of insert-row events.

- If the statement fails, nothing is logged. This includes the case that the destination table exists and `IF NOT EXISTS` is not given.

When the destination table exists and `IF NOT EXISTS` is given, MySQL 8.0 ignores the statement completely; nothing is inserted or logged. The handling of such statements in this regard has changed considerably in previous MySQL releases; if you are replicating from a MySQL 5.5.6 or older master to a newer slave, see Replication of CREATE ... IF NOT EXISTS Statements, for more information.

18.4.1.7 Replication of CREATE SERVER, ALTER SERVER, and DROP SERVER

In MySQL 8.0, the statements `CREATE SERVER`, `ALTER SERVER`, and `DROP SERVER` are not written to the binary log, regardless of the binary logging format that is in use.

18.4.1.8 Replication of CURRENT_USER()

The following statements support use of the `CURRENT_USER()` function to take the place of the name of (and, possibly, the host for) an affected user or a definer; in such cases, `CURRENT_USER()` is expanded where and as needed:

- `DROP USER`
- `RENAME USER`
- `GRANT`
- `REVOKE`
- `CREATE FUNCTION`
- `CREATE PROCEDURE`
- `CREATE TRIGGER`
- `CREATE EVENT`
- `CREATE VIEW`
- `ALTER EVENT`
- `ALTER VIEW`
- `SET PASSWORD`

When `CURRENT_USER()` or `CURRENT_USER` is used as the definer in any of the statements `CREATE FUNCTION`, `CREATE PROCEDURE`, `CREATE TRIGGER`, `CREATE EVENT`, `CREATE VIEW`, or `ALTER VIEW` when binary logging is enabled, the function reference is expanded before it is written to the binary log, so that the statement refers to the same user on both the master and the slave when the statement is replicated. `CURRENT_USER()` or `CURRENT_USER` is also expanded prior to being written to the binary log when used in `DROP USER`, `RENAME USER`, `GRANT`, `REVOKE`, or `ALTER EVENT`.

18.4.1.9 Replication of DROP ... IF EXISTS Statements

The `DROP DATABASE IF EXISTS`, `DROP TABLE IF EXISTS`, and `DROP VIEW IF EXISTS` statements are always replicated, even if the database, table, or view to be dropped does not exist on the master. This is to ensure that the object to be dropped no longer exists on either the master or the slave, once the slave has caught up with the master.

`DROP ... IF EXISTS` statements for stored programs (stored procedures and functions, triggers, and events) are also replicated, even if the stored program to be dropped does not exist on the master.

18.4.1.10 Replication with Differing Table Definitions on Master and Slave

Source and target tables for replication do not have to be identical. A table on the master can have more or fewer columns than the slave's copy of the table. In addition, corresponding table columns on the master and the slave can use different data types, subject to certain conditions.

Note

Replication between tables which are partitioned differently from one another is not supported. See Section 18.4.1.19, “Replication and Partitioning”.

In all cases where the source and target tables do not have identical definitions, the database and table names must be the same on both the master and the slave. Additional conditions are discussed, with examples, in the following two sections.

Replication with More Columns on Master or Slave

You can replicate a table from the master to the slave such that the master and slave copies of the table have differing numbers of columns, subject to the following conditions:

- Columns common to both versions of the table must be defined in the same order on the master and the slave.

  (This is true even if both tables have the same number of columns.)

- Columns common to both versions of the table must be defined before any additional columns.

  This means that executing an `ALTER TABLE` statement on the slave where a new column is inserted into the table within the range of columns common to both tables causes replication to fail, as shown in the following example:

  Suppose that a table `t`, existing on the master and the slave, is defined by the following `CREATE TABLE` statement:

  ```
  CREATE TABLE t (
      c1 INT,
      c2 INT,
      c3 INT
  );
  ```

  Suppose that the `ALTER TABLE` statement shown here is executed on the slave:

  ```
  ALTER TABLE t ADD COLUMN cnew1 INT AFTER c3;
  ```

  The previous `ALTER TABLE` is permitted on the slave because the columns `c1`, `c2`, and `c3` that are common to both versions of table `t` remain grouped together in both versions of the table, before any columns that differ.

  However, the following `ALTER TABLE` statement cannot be executed on the slave without causing replication to break:

  ```
  ALTER TABLE t ADD COLUMN cnew2 INT AFTER c2;
  ```

  Replication fails after execution on the slave of the `ALTER TABLE` statement just shown, because the new column `cnew2` comes between columns common to both versions of `t`.

- Each “extra” column in the version of the table having more columns must have a default value.

  A column's default value is determined by a number of factors, including its type, whether it is defined with a `DEFAULT` option, whether it is declared as `NULL`, and the server SQL mode in effect at the time of its creation; for more information, see Section 11.7, “Data Type Default Values”).

In addition, when the slave's copy of the table has more columns than the master's copy, each column common to the tables must use the same data type in both tables.

Examples. The following examples illustrate some valid and invalid table definitions:

More columns on the master. The following table definitions are valid and replicate correctly:

```
master> CREATE TABLE t1 (c1 INT, c2 INT, c3 INT);
slave>  CREATE TABLE t1 (c1 INT, c2 INT);
```

The following table definitions would raise an error because the definitions of the columns common to both versions of the table are in a different order on the slave than they are on the master:

```
master> CREATE TABLE t1 (c1 INT, c2 INT, c3 INT);
slave>  CREATE TABLE t1 (c2 INT, c1 INT);
```

The following table definitions would also raise an error because the definition of the extra column on the master appears before the definitions of the columns common to both versions of the table:

```
master> CREATE TABLE t1 (c3 INT, c1 INT, c2 INT);
slave>  CREATE TABLE t1 (c1 INT, c2 INT);
```

More columns on the slave. The following table definitions are valid and replicate correctly:

```
master> CREATE TABLE t1 (c1 INT, c2 INT);
slave>  CREATE TABLE t1 (c1 INT, c2 INT, c3 INT);
```

The following definitions raise an error because the columns common to both versions of the table are not defined in the same order on both the master and the slave:

```
master> CREATE TABLE t1 (c1 INT, c2 INT);
slave>  CREATE TABLE t1 (c2 INT, c1 INT, c3 INT);
```

The following table definitions also raise an error because the definition for the extra column in the slave's version of the table appears before the definitions for the columns which are common to both versions of the table:

```
master> CREATE TABLE t1 (c1 INT, c2 INT);
slave>  CREATE TABLE t1 (c3 INT, c1 INT, c2 INT);
```

The following table definitions fail because the slave's version of the table has additional columns compared to the master's version, and the two versions of the table use different data types for the common column `c2`:

```
master> CREATE TABLE t1 (c1 INT, c2 BIGINT);
slave>  CREATE TABLE t1 (c1 INT, c2 INT, c3 INT);
```

Replication of Columns Having Different Data Types

Corresponding columns on the master's and the slave's copies of the same table ideally should have the same data type. However, this is not always strictly enforced, as long as certain conditions are met.

It is usually possible to replicate from a column of a given data type to another column of the same type and same size or width, where applicable, or larger. For example, you can replicate from a `CHAR(10)` column to another `CHAR(10)`, or from a `CHAR(10)` column to a `CHAR(25)` column without any problems. In certain cases, it also possible to replicate from a column having one data type (on

the master) to a column having a different data type (on the slave); when the data type of the master's version of the column is promoted to a type that is the same size or larger on the slave, this is known as *attribute promotion*.

Attribute promotion can be used with both statement-based and row-based replication, and is not dependent on the storage engine used by either the master or the slave. However, the choice of logging format does have an effect on the type conversions that are permitted; the particulars are discussed later in this section.

Important

Whether you use statement-based or row-based replication, the slave's copy of the table cannot contain more columns than the master's copy if you wish to employ attribute promotion.

Statement-based replication. When using statement-based replication, a simple rule of thumb to follow is, “If the statement run on the master would also execute successfully on the slave, it should also replicate successfully”. In other words, if the statement uses a value that is compatible with the type of a given column on the slave, the statement can be replicated. For example, you can insert any value that fits in a `TINYINT` column into a `BIGINT` column as well; it follows that, even if you change the type of a `TINYINT` column in the slave's copy of a table to `BIGINT`, any insert into that column on the master that succeeds should also succeed on the slave, since it is impossible to have a legal `TINYINT` value that is large enough to exceed a `BIGINT` column.

Row-based replication: attribute promotion and demotion. Row-based replication in MySQL 8.0 supports attribute promotion and demotion between smaller data types and larger types. It is also possible to specify whether or not to permit lossy (truncated) or non-lossy conversions of demoted column values, as explained later in this section.

Lossy and non-lossy conversions. In the event that the target type cannot represent the value being inserted, a decision must be made on how to handle the conversion. If we permit the conversion but truncate (or otherwise modify) the source value to achieve a “fit” in the target column, we make what is known as a *lossy conversion*. A conversion which does not require truncation or similar modifications to fit the source column value in the target column is a *non-lossy* conversion.

Type conversion modes (slave_type_conversions variable). The setting of the `slave_type_conversions` global server variable controls the type conversion mode used on the slave. This variable takes a set of values from the following table, which shows the effects of each mode on the slave's type-conversion behavior:

Mode	Effect
`ALL_LOSSY`	In this mode, type conversions that would mean loss of information are permitted. This does not imply that non-lossy conversions are permitted, merely that only cases requiring either lossy conversions or no conversion at all are permitted; for example, enabling *only* this mode permits an `INT` column to be converted to `TINYINT` (a lossy conversion), but not a `TINYINT` column to an `INT` column (non-lossy). Attempting the latter conversion in this case would cause replication to stop with an error on the slave.
`ALL_NON_LOSSY`	This mode permits conversions that do not require truncation or other special handling of the source value; that is, it permits conversions where the target type has a wider range than the source type. Setting this mode has no bearing on whether lossy conversions are permitted; this is controlled with the `ALL_LOSSY` mode. If only

Mode	Effect
	ALL_NON_LOSSY is set, but not ALL_LOSSY, then attempting a conversion that would result in the loss of data (such as INT to TINYINT, or CHAR(25) to VARCHAR(20)) causes the slave to stop with an error.
ALL_LOSSY,ALL_NON_LOSSY	When this mode is set, all supported type conversions are permitted, whether or not they are lossy conversions.
ALL_SIGNED	Treat promoted integer types as signed values (the default behavior).
ALL_UNSIGNED	Treat promoted integer types as unsigned values.
ALL_SIGNED,ALL_UNSIGNED	Treat promoted integer types as signed if possible, otherwise as unsigned.
[*empty*]	When slave_type_conversions is not set, no attribute promotion or demotion is permitted; this means that all columns in the source and target tables must be of the same types. This mode is the default.

When an integer type is promoted, its signedness is not preserved. By default, the slave treats all such values as signed. You can control this behavior using ALL_SIGNED, ALL_UNSIGNED, or both. ALL_SIGNED tells the slave to treat all promoted integer types as signed; ALL_UNSIGNED instructs it to treat these as unsigned. Specifying both causes the slave to treat the value as signed if possible, otherwise to treat it as unsigned; the order in which they are listed is not significant. Neither ALL_SIGNED nor ALL_UNSIGNED has any effect if at least one of ALL_LOSSY or ALL_NONLOSSY is not also used.

Changing the type conversion mode requires restarting the slave with the new slave_type_conversions setting.

Supported conversions. Supported conversions between different but similar data types are shown in the following list:

- Between any of the integer types TINYINT, SMALLINT, MEDIUMINT, INT, and BIGINT.

  This includes conversions between the signed and unsigned versions of these types.

  Lossy conversions are made by truncating the source value to the maximum (or minimum) permitted by the target column. For ensuring non-lossy conversions when going from unsigned to signed types, the target column must be large enough to accommodate the range of values in the source column. For example, you can demote TINYINT UNSIGNED non-lossily to SMALLINT, but not to TINYINT.

- Between any of the decimal types DECIMAL, FLOAT, DOUBLE, and NUMERIC.

  FLOAT to DOUBLE is a non-lossy conversion; DOUBLE to FLOAT can only be handled lossily. A conversion from DECIMAL(*M*,*D*) to DECIMAL(*M'*,*D'*) where *D'* >= *D* and (*M'*-*D'*) >= (*M*-*D*) is non-lossy; for any case where *M'* < *M*, *D'* < *D*, or both, only a lossy conversion can be made.

  For any of the decimal types, if a value to be stored cannot be fit in the target type, the value is rounded down according to the rounding rules defined for the server elsewhere in the documentation. See Section 12.21.4, “Rounding Behavior”, for information about how this is done for decimal types.

- Between any of the string types CHAR, VARCHAR, and TEXT, including conversions between different widths.

  Conversion of a CHAR, VARCHAR, or TEXT to a CHAR, VARCHAR, or TEXT column the same size or larger is never lossy. Lossy conversion is handled by inserting only the first *N* characters of the string on the slave, where *N* is the width of the target column.

Important

Replication between columns using different character sets is not supported.

- Between any of the binary data types `BINARY`, `VARBINARY`, and `BLOB`, including conversions between different widths.

  Conversion of a `BINARY`, `VARBINARY`, or `BLOB` to a `BINARY`, `VARBINARY`, or `BLOB` column the same size or larger is never lossy. Lossy conversion is handled by inserting only the first *`N`* bytes of the string on the slave, where *`N`* is the width of the target column.

- Between any 2 `BIT` columns of any 2 sizes.

  When inserting a value from a `BIT(`*`M`*`)` column into a `BIT(`*`M'`*`)` column, where *`M' > M`*, the most significant bits of the `BIT(`*`M'`*`)` columns are cleared (set to zero) and the *`M`* bits of the `BIT(`*`M`*`)` value are set as the least significant bits of the `BIT(`*`M'`*`)` column.

  When inserting a value from a source `BIT(`*`M`*`)` column into a target `BIT(`*`M'`*`)` column, where *`M' < M`*, the maximum possible value for the `BIT(`*`M'`*`)` column is assigned; in other words, an “all-set” value is assigned to the target column.

Conversions between types not in the previous list are not permitted.

18.4.1.11 Replication and DIRECTORY Table Options

If a `DATA DIRECTORY` or `INDEX DIRECTORY` table option is used in a `CREATE TABLE` statement on the master server, the table option is also used on the slave. This can cause problems if no corresponding directory exists in the slave host file system or if it exists but is not accessible to the slave server. This can be overridden by using the `NO_DIR_IN_CREATE` server SQL mode on the slave, which causes the slave to ignore the `DATA DIRECTORY` and `INDEX DIRECTORY` table options when replicating `CREATE TABLE` statements. The result is that `MyISAM` data and index files are created in the table's database directory.

For more information, see Section 5.1.8, “Server SQL Modes”.

18.4.1.12 Replication of Invoked Features

Replication of invoked features such as user-defined functions (UDFs) and stored programs (stored procedures and functions, triggers, and events) provides the following characteristics:

- The effects of the feature are always replicated.
- The following statements are replicated using statement-based replication:
  - `CREATE EVENT`
  - `ALTER EVENT`
  - `DROP EVENT`
  - `CREATE PROCEDURE`
  - `DROP PROCEDURE`
  - `CREATE FUNCTION`
  - `DROP FUNCTION`
  - `CREATE TRIGGER`
  - `DROP TRIGGER`

However, the *effects* of features created, modified, or dropped using these statements are replicated using row-based replication.

Note

Attempting to replicate invoked features using statement-based replication produces the warning `Statement is not safe to log in statement format`. For example, trying to replicate a UDF with statement-based replication generates this warning because it currently cannot be determined by the MySQL server whether the UDF is deterministic. If you are absolutely certain that the invoked feature's effects are deterministic, you can safely disregard such warnings.

- In the case of `CREATE EVENT` and `ALTER EVENT`:
  - The status of the event is set to `SLAVESIDE_DISABLED` on the slave regardless of the state specified (this does not apply to `DROP EVENT`).
  - The master on which the event was created is identified on the slave by its server ID. The `ORIGINATOR` column in `INFORMATION_SCHEMA.EVENTS` and the `originator` column in `mysql.event` store this information. See Section 24.8, "The INFORMATION_SCHEMA EVENTS Table", and Section 13.7.5.18, "SHOW EVENTS Syntax", for more information.
- The feature implementation resides on the slave in a renewable state so that if the master fails, the slave can be used as the master without loss of event processing.

To determine whether there are any scheduled events on a MySQL server that were created on a different server (that was acting as a replication master), query the `INFORMATION_SCHEMA.EVENTS` table in a manner similar to what is shown here:

```
SELECT EVENT_SCHEMA, EVENT_NAME
    FROM INFORMATION_SCHEMA.EVENTS
    WHERE STATUS = 'SLAVESIDE_DISABLED';
```

Alternatively, you can use the `SHOW EVENTS` statement, like this:

```
SHOW EVENTS
    WHERE STATUS = 'SLAVESIDE_DISABLED';
```

When promoting a replication slave having such events to a replication master, you must enable each event using `ALTER EVENT event_name ENABLE`, where *`event_name`* is the name of the event.

If more than one master was involved in creating events on this slave, and you wish to identify events that were created only on a given master having the server ID *`master_id`*, modify the previous query on the `EVENTS` table to include the `ORIGINATOR` column, as shown here:

```
SELECT EVENT_SCHEMA, EVENT_NAME, ORIGINATOR
    FROM INFORMATION_SCHEMA.EVENTS
    WHERE STATUS = 'SLAVESIDE_DISABLED'
    AND   ORIGINATOR = 'master_id'
```

You can employ `ORIGINATOR` with the `SHOW EVENTS` statement in a similar fashion:

```
SHOW EVENTS
    WHERE STATUS = 'SLAVESIDE_DISABLED'
    AND   ORIGINATOR = 'master_id'
```

Before enabling events that were replicated from the master, you should disable the MySQL Event Scheduler on the slave (using a statement such as `SET GLOBAL event_scheduler = OFF;`), run

any necessary ALTER EVENT statements, restart the server, then re-enable the Event Scheduler on the slave afterward (using a statement such as SET GLOBAL event_scheduler = ON;)-

If you later demote the new master back to being a replication slave, you must disable manually all events enabled by the ALTER EVENT statements. You can do this by storing in a separate table the event names from the SELECT statement shown previously, or using ALTER EVENT statements to rename the events with a common prefix such as replicated_ to identify them.

If you rename the events, then when demoting this server back to being a replication slave, you can identify the events by querying the EVENTS table, as shown here:

```
SELECT CONCAT(EVENT_SCHEMA, '.', EVENT_NAME) AS 'Db.Event'
      FROM INFORMATION_SCHEMA.EVENTS
      WHERE INSTR(EVENT_NAME, 'replicated_') = 1;
```

18.4.1.13 Replication and Floating-Point Values

With statement-based replication, values are converted from decimal to binary. Because conversions between decimal and binary representations of them may be approximate, comparisons involving floating-point values are inexact. This is true for operations that use floating-point values explicitly, or that use values that are converted to floating-point implicitly. Comparisons of floating-point values might yield different results on master and slave servers due to differences in computer architecture, the compiler used to build MySQL, and so forth. See Section 12.2, “Type Conversion in Expression Evaluation”, and Section B.5.4.8, “Problems with Floating-Point Values”.

18.4.1.14 Replication and Fractional Seconds Support

MySQL 8.0 permits fractional seconds for TIME, DATETIME, and TIMESTAMP values, with up to microseconds (6 digits) precision. See Section 11.3.6, “Fractional Seconds in Time Values”.

There may be problems replicating from a master server that understands fractional seconds to an older slave (MySQL 5.6.3 and earlier) that does not:

- For CREATE TABLE statements containing columns that have an *fsp* (fractional seconds precision) value greater than 0, replication will fail due to parser errors.
- Statements that use temporal data types with an *fsp* value of 0 will work for with statement-based logging but not row-based logging. In the latter case, the data types have binary formats and type codes on the master that differ from those on the slave.
- Some expression results will differ on master and slave. Examples: On the master, the timestamp system variable returns a value that includes a microseconds fractional part; on the slave, it returns an integer. On the master, functions that return a result that includes the current time (such as CURTIME(), SYSDATE(), or UTC_TIMESTAMP()) interpret an argument as an *fsp* value and the return value includes a fractional seconds part of that many digits. On the slave, these functions permit an argument but ignore it.

18.4.1.15 Replication and FLUSH

Some forms of the FLUSH statement are not logged because they could cause problems if replicated to a slave: FLUSH LOGS, FLUSH MASTER, FLUSH SLAVE, and FLUSH TABLES WITH READ LOCK. For a syntax example, see Section 13.7.6.3, “FLUSH Syntax”. The FLUSH TABLES, ANALYZE TABLE, OPTIMIZE TABLE, and REPAIR TABLE statements are written to the binary log and thus replicated to slaves. This is not normally a problem because these statements do not modify table data.

However, this behavior can cause difficulties under certain circumstances. If you replicate the privilege tables in the mysql database and update those tables directly without using GRANT, you must issue a FLUSH PRIVILEGES on the slaves to put the new privileges into effect. In addition, if you use FLUSH TABLES when renaming a MyISAM table that is part of a MERGE table, you must issue FLUSH

TABLES manually on the slaves. These statements are written to the binary log unless you specify NO_WRITE_TO_BINLOG or its alias LOCAL.

18.4.1.16 Replication and System Functions

Certain functions do not replicate well under some conditions:

- The USER(), CURRENT_USER() (or CURRENT_USER), UUID(), VERSION(), and LOAD_FILE() functions are replicated without change and thus do not work reliably on the slave unless row-based replication is enabled. (See Section 18.2.1, “Replication Formats”.)

  USER() and CURRENT_USER() are automatically replicated using row-based replication when using MIXED mode, and generate a warning in STATEMENT mode. (See also Section 18.4.1.8, “Replication of CURRENT_USER()”.) This is also true for VERSION() and RAND().

- For NOW(), the binary log includes the timestamp. This means that the value *as returned by the call to this function on the master* is replicated to the slave. To avoid unexpected results when replicating between MySQL servers in different time zones, set the time zone on both master and slave. See also Section 18.4.1.32, “Replication and Time Zones”

  To explain the potential problems when replicating between servers which are in different time zones, suppose that the master is located in New York, the slave is located in Stockholm, and both servers are using local time. Suppose further that, on the master, you create a table mytable, perform an INSERT statement on this table, and then select from the table, as shown here:

```
mysql> CREATE TABLE mytable (mycol TEXT);
Query OK, 0 rows affected (0.06 sec)

mysql> INSERT INTO mytable VALUES ( NOW() );
Query OK, 1 row affected (0.00 sec)

mysql> SELECT * FROM mytable;
+---------------------+
| mycol               |
+---------------------+
| 2009-09-01 12:00:00 |
+---------------------+
1 row in set (0.00 sec)
```

  Local time in Stockholm is 6 hours later than in New York; so, if you issue SELECT NOW() on the slave at that exact same instant, the value 2009-09-01 18:00:00 is returned. For this reason, if you select from the slave's copy of mytable after the CREATE TABLE and INSERT statements just shown have been replicated, you might expect mycol to contain the value 2009-09-01 18:00:00. However, this is not the case; when you select from the slave's copy of mytable, you obtain exactly the same result as on the master:

```
mysql> SELECT * FROM mytable;
+---------------------+
| mycol               |
+---------------------+
| 2009-09-01 12:00:00 |
+---------------------+
1 row in set (0.00 sec)
```

  Unlike NOW(), the SYSDATE() function is not replication-safe because it is not affected by SET TIMESTAMP statements in the binary log and is nondeterministic if statement-based logging is used. This is not a problem if row-based logging is used.

  An alternative is to use the --sysdate-is-now option to cause SYSDATE() to be an alias for NOW(). This must be done on the master and the slave to work correctly. In such cases, a warning is still issued by this function, but can safely be ignored as long as --sysdate-is-now is used on both the master and the slave.

`SYSDATE()` is automatically replicated using row-based replication when using `MIXED` mode, and generates a warning in `STATEMENT` mode.

See also Section 18.4.1.32, “Replication and Time Zones”.

- *The following restriction applies to statement-based replication only, not to row-based replication.* The `GET_LOCK()`, `RELEASE_LOCK()`, `IS_FREE_LOCK()`, and `IS_USED_LOCK()` functions that handle user-level locks are replicated without the slave knowing the concurrency context on the master. Therefore, these functions should not be used to insert into a master table because the content on the slave would differ. For example, do not issue a statement such as `INSERT INTO mytable VALUES(GET_LOCK(...))`.

  These functions are automatically replicated using row-based replication when using `MIXED` mode, and generate a warning in `STATEMENT` mode.

As a workaround for the preceding limitations when statement-based replication is in effect, you can use the strategy of saving the problematic function result in a user variable and referring to the variable in a later statement. For example, the following single-row `INSERT` is problematic due to the reference to the `UUID()` function:

```
INSERT INTO t VALUES(UUID());
```

To work around the problem, do this instead:

```
SET @my_uuid = UUID();
INSERT INTO t VALUES(@my_uuid);
```

That sequence of statements replicates because the value of `@my_uuid` is stored in the binary log as a user-variable event prior to the `INSERT` statement and is available for use in the `INSERT`.

The same idea applies to multiple-row inserts, but is more cumbersome to use. For a two-row insert, you can do this:

```
SET @my_uuid1 = UUID(); @my_uuid2 = UUID();
INSERT INTO t VALUES(@my_uuid1),(@my_uuid2);
```

However, if the number of rows is large or unknown, the workaround is difficult or impracticable. For example, you cannot convert the following statement to one in which a given individual user variable is associated with each row:

```
INSERT INTO t2 SELECT UUID(), * FROM t1;
```

Within a stored function, `RAND()` replicates correctly as long as it is invoked only once during the execution of the function. (You can consider the function execution timestamp and random number seed as implicit inputs that are identical on the master and slave.)

The `FOUND_ROWS()` and `ROW_COUNT()` functions are not replicated reliably using statement-based replication. A workaround is to store the result of the function call in a user variable, and then use that in the `INSERT` statement. For example, if you wish to store the result in a table named `mytable`, you might normally do so like this:

```
SELECT SQL_CALC_FOUND_ROWS FROM mytable LIMIT 1;
INSERT INTO mytable VALUES( FOUND_ROWS() );
```

However, if you are replicating `mytable`, you should use `SELECT ... INTO`, and then store the variable in the table, like this:

```
SELECT SQL_CALC_FOUND_ROWS INTO @found_rows FROM mytable LIMIT 1;
INSERT INTO mytable VALUES(@found_rows);
```

In this way, the user variable is replicated as part of the context, and applied on the slave correctly.

These functions are automatically replicated using row-based replication when using `MIXED` mode, and generate a warning in `STATEMENT` mode. (Bug #12092, Bug #30244)

18.4.1.17 Replication and LIMIT

Statement-based replication of `LIMIT` clauses in `DELETE`, `UPDATE`, and `INSERT ... SELECT` statements is unsafe since the order of the rows affected is not defined. (Such statements can be replicated correctly with statement-based replication only if they also contain an `ORDER BY` clause.) When such a statement is encountered:

- When using `STATEMENT` mode, a warning that the statement is not safe for statement-based replication is now issued.

  When using `STATEMENT` mode, warnings are issued for DML statements containing `LIMIT` even when they also have an `ORDER BY` clause (and so are made deterministic). This is a known issue. (Bug #42851)

- When using `MIXED` mode, the statement is now automatically replicated using row-based mode.

18.4.1.18 Replication and LOAD DATA INFILE

In MySQL 8.0, `LOAD DATA INFILE` is considered unsafe (see Section 18.2.1.3, “Determination of Safe and Unsafe Statements in Binary Logging”).

18.4.1.19 Replication and Partitioning

Replication is supported between partitioned tables as long as they use the same partitioning scheme and otherwise have the same structure except where an exception is specifically allowed (see Section 18.4.1.10, “Replication with Differing Table Definitions on Master and Slave”).

Replication between tables having different partitioning is generally not supported. This because statements (such as `ALTER TABLE ... DROP PARTITION`) acting directly on partitions in such cases may produce different results on master and slave. In the case where a table is partitioned on the master but not on the slave, any statements operating on partitions on the master's copy of the slave fail on the slave. When the slave's copy of the table is partitioned but the master's copy is not, statements acting on partitions cannot be run on the master without causing errors there.

Due to these dangers of causing replication to fail entirely (on account of failed statements) and of inconsistencies (when the result of a partition-level SQL statement produces different results on master and slave), we recommend that insure that the partitioning of any tables to be replicated from the master is matched by the slave's versions of these tables.

18.4.1.20 Replication and REPAIR TABLE

When used on a corrupted or otherwise damaged table, it is possible for the `REPAIR TABLE` statement to delete rows that cannot be recovered. However, any such modifications of table data performed by this statement are not replicated, which can cause master and slave to lose synchronization. For this reason, in the event that a table on the master becomes damaged and you use `REPAIR TABLE` to repair it, you should first stop replication (if it is still running) before using `REPAIR TABLE`, then afterward compare the master's and slave's copies of the table and be prepared to correct any discrepancies manually, before restarting replication.

18.4.1.21 Replication and Master or Slave Shutdowns

It is safe to shut down a master server and restart it later. When a slave loses its connection to the master, the slave tries to reconnect immediately and retries periodically if that fails. The default is to retry every 60 seconds. This may be changed with the `CHANGE MASTER TO` statement. A slave also is able to deal with network connectivity outages. However, the slave notices the network outage only after receiving no data from the master for `slave_net_timeout` seconds. If your outages are short,

you may want to decrease `slave_net_timeout`. See Section 18.3.2, “Handling an Unexpected Halt of a Replication Slave”.

An unclean shutdown (for example, a crash) on the master side can result in the master binary log having a final position less than the most recent position read by the slave, due to the master binary log file not being flushed. This can cause the slave not to be able to replicate when the master comes back up. Setting `sync_binlog=1` in the master `my.cnf` file helps to minimize this problem because it causes the master to flush its binary log more frequently.

Shutting down a slave cleanly is safe because it keeps track of where it left off. However, be careful that the slave does not have temporary tables open; see Section 18.4.1.24, “Replication and Temporary Tables”. Unclean shutdowns might produce problems, especially if the disk cache was not flushed to disk before the problem occurred:

- For transactions, the slave commits and then updates `relay-log.info`. If a crash occurs between these two operations, relay log processing will have proceeded further than the information file indicates and the slave will re-execute the events from the last transaction in the relay log after it has been restarted.
- A similar problem can occur if the slave updates `relay-log.info` but the server host crashes before the write has been flushed to disk. To minimize the chance of this occurring, set `sync_relay_log_info=1` in the slave `my.cnf` file. Setting `sync_relay_log_info` to 0 causes no writes to be forced to disk and the server relies on the operating system to flush the file from time to time.

The fault tolerance of your system for these types of problems is greatly increased if you have a good uninterruptible power supply.

18.4.1.22 Replication and max_allowed_packet

`max_allowed_packet` sets an upper limit on the size of any single message between the MySQL server and clients, including replication slaves. If you are replicating large column values (such as might be found in `TEXT` or `BLOB` columns) and `max_allowed_packet` is too small on the master, the master fails with an error, and the slave shuts down the I/O thread. If `max_allowed_packet` is too small on the slave, this also causes the slave to stop the I/O thread.

Row-based replication currently sends all columns and column values for updated rows from the master to the slave, including values of columns that were not actually changed by the update. This means that, when you are replicating large column values using row-based replication, you must take care to set `max_allowed_packet` large enough to accommodate the largest row in any table to be replicated, even if you are replicating updates only, or you are inserting only relatively small values.

18.4.1.23 Replication and MEMORY Tables

When a master server shuts down and restarts, its `MEMORY` tables become empty. To replicate this effect to slaves, the first time that the master uses a given `MEMORY` table after startup, it logs an event that notifies slaves that the table must to be emptied by writing a `DELETE` statement for that table to the binary log.

When a slave server shuts down and restarts, its `MEMORY` tables become empty. This causes the slave to be out of synchrony with the master and may lead to other failures or cause the slave to stop:

- Row-format updates and deletes received from the master may fail with `Can't find record in 'memory_table'`.
- Statements such as `INSERT INTO ... SELECT FROM memory_table` may insert a different set of rows on the master and slave.

The safe way to restart a slave that is replicating `MEMORY` tables is to first drop or delete all rows from the `MEMORY` tables on the master and wait until those changes have replicated to the slave. Then it is safe to restart the slave.

An alternative restart method may apply in some cases. When `binlog_format=ROW`, you can prevent the slave from stopping if you set `slave_exec_mode=IDEMPOTENT` before you start the slave again. This allows the slave to continue to replicate, but its `MEMORY` tables will still be different from those on the master. This can be okay if the application logic is such that the contents of `MEMORY` tables can be safely lost (for example, if the `MEMORY` tables are used for caching). `slave_exec_mode=IDEMPOTENT` applies globally to all tables, so it may hide other replication errors in non-`MEMORY` tables.

The size of `MEMORY` tables is limited by the value of the `max_heap_table_size` system variable, which is not replicated (see Section 18.4.1.38, "Replication and Variables"). A change in `max_heap_table_size` takes effect for `MEMORY` tables that are created or updated using `ALTER TABLE ... ENGINE = MEMORY` or `TRUNCATE TABLE` following the change, or for all `MEMORY` tables following a server restart. If you increase the value of this variable on the master without doing so on the slave, it becomes possible for a table on the master to grow larger than its counterpart on the slave, leading to inserts that succeed on the master but fail on the slave with `Table is full` errors. This is a known issue (Bug #48666). In such cases, you must set the global value of `max_heap_table_size` on the slave as well as on the master, then restart replication. It is also recommended that you restart both the master and slave MySQL servers, to insure that the new value takes complete (global) effect on each of them.

See Section 16.3, "The MEMORY Storage Engine", for more information about `MEMORY` tables.

18.4.1.24 Replication and Temporary Tables

The discussion in the following paragraphs does not apply when `binlog_format=ROW` because, in that case, temporary tables are not replicated; this means that there are never any temporary tables on the slave to be lost in the event of an unplanned shutdown by the slave. The remainder of this section applies only when using statement-based or mixed-format replication. Loss of replicated temporary tables on the slave can be an issue, whenever `binlog_format` is `STATEMENT` or `MIXED`, for statements involving temporary tables that can be logged safely using statement-based format. For more information about row-based replication and temporary tables, see Row-based logging of temporary tables.

Safe slave shutdown when using temporary tables. Temporary tables are replicated except in the case where you stop the slave server (not just the slave threads) and you have replicated temporary tables that are open for use in updates that have not yet been executed on the slave. If you stop the slave server, the temporary tables needed by those updates are no longer available when the slave is restarted. To avoid this problem, do not shut down the slave while it has temporary tables open. Instead, use the following procedure:

1. Issue a `STOP SLAVE SQL_THREAD` statement.
2. Use `SHOW STATUS` to check the value of the `Slave_open_temp_tables` variable.
3. If the value is not 0, restart the slave SQL thread with `START SLAVE SQL_THREAD` and repeat the procedure later.
4. When the value is 0, issue a `mysqladmin shutdown` command to stop the slave.

Temporary tables and replication options. By default, all temporary tables are replicated; this happens whether or not there are any matching `--replicate-do-db`, `--replicate-do-table`, or `--replicate-wild-do-table` options in effect. However, the `--replicate-ignore-table` and `--replicate-wild-ignore-table` options are honored for temporary tables.

A recommended practice when using statement-based or mixed-format replication is to designate a prefix for exclusive use in naming temporary tables that you do not want replicated, then employ a `--replicate-wild-ignore-table` option to match that prefix. For example, you might give all such tables names beginning with `norep` (such as `norepmytable`, `norepyourtable`, and so on), then use `--replicate-wild-ignore-table=norep%` to prevent them from being replicated.

18.4.1.25 Replication of the mysql System Database

Data modification statements made to tables in the `mysql` database are replicated according to the value of `binlog_format`; if this value is `MIXED`, these statements are replicated using row-based format. However, statements that would normally update this information indirectly—such `GRANT`, `REVOKE`, and statements manipulating triggers, stored routines, and views—are replicated to slaves using statement-based replication.

18.4.1.26 Replication and the Query Optimizer

It is possible for the data on the master and slave to become different if a statement is written in such a way that the data modification is nondeterministic; that is, left up the query optimizer. (In general, this is not a good practice, even outside of replication.) Examples of nondeterministic statements include `DELETE` or `UPDATE` statements that use `LIMIT` with no `ORDER BY` clause; see Section 18.4.1.17, “Replication and LIMIT”, for a detailed discussion of these.

18.4.1.27 Replication and Reserved Words

You can encounter problems when you attempt to replicate from an older master to a newer slave and you make use of identifiers on the master that are reserved words in the newer MySQL version running on the slave. An example of this is using a table column named `virtual` on a 5.6 master that is replicating to a 5.7 or higher slave because `VIRTUAL` is a reserved word beginning in MySQL 5.7. Replication can fail in such cases with Error 1064 `You have an error in your SQL syntax...`, *even if a database or table named using the reserved word or a table having a column named using the reserved word is excluded from replication*. This is due to the fact that each SQL event must be parsed by the slave prior to execution, so that the slave knows which database object or objects would be affected; only after the event is parsed can the slave apply any filtering rules defined by `--replicate-do-db`, `--replicate-do-table`, `--replicate-ignore-db`, and `--replicate-ignore-table`.

To work around the problem of database, table, or column names on the master which would be regarded as reserved words by the slave, do one of the following:

- Use one or more `ALTER TABLE` statements on the master to change the names of any database objects where these names would be considered reserved words on the slave, and change any SQL statements that use the old names to use the new names instead.
- In any SQL statements using these database object names, write the names as quoted identifiers using backtick characters (`` ` ``).

For listings of reserved words by MySQL version, see Reserved Words, in the *MySQL Server Version Reference*. For identifier quoting rules, see Section 9.2, “Schema Object Names”.

18.4.1.28 Slave Errors During Replication

If a statement produces the same error (identical error code) on both the master and the slave, the error is logged, but replication continues.

If a statement produces different errors on the master and the slave, the slave SQL thread terminates, and the slave writes a message to its error log and waits for the database administrator to decide what to do about the error. This includes the case that a statement produces an error on the master or the slave, but not both. To address the issue, connect to the slave manually and determine the cause of the problem. `SHOW SLAVE STATUS` is useful for this. Then fix the problem and run `START SLAVE`. For example, you might need to create a nonexistent table before you can start the slave again.

If this error code validation behavior is not desirable, some or all errors can be masked out (ignored) with the `--slave-skip-errors` option.

For nontransactional storage engines such as `MyISAM`, it is possible to have a statement that only partially updates a table and returns an error code. This can happen, for example, on a multiple-row insert that has one row violating a key constraint, or if a long update statement is killed after updating

some of the rows. If that happens on the master, the slave expects execution of the statement to result in the same error code. If it does not, the slave SQL thread stops as described previously.

If you are replicating between tables that use different storage engines on the master and slave, keep in mind that the same statement might produce a different error when run against one version of the table, but not the other, or might cause an error for one version of the table, but not the other. For example, since `MyISAM` ignores foreign key constraints, an `INSERT` or `UPDATE` statement accessing an `InnoDB` table on the master might cause a foreign key violation but the same statement performed on a `MyISAM` version of the same table on the slave would produce no such error, causing replication to stop.

18.4.1.29 Replication of Server-Side Help Tables

The server maintains tables in the `mysql` database that store information for the `HELP` statement (see Section 13.8.3, “HELP Syntax”. These tables can be loaded manually as described at Section 5.1.10, “Server-Side Help”.

Help table content is derived from the MySQL Reference Manual. There are versions of the manual specific to each MySQL release series, so help content is specific to each series as well. Normally, you load a version of help content that matches the server version. This has implications for replication. For example, you would load MySQL 5.6 help content into a MySQL 5.6 master server, but not necessarily replicate that content to a MySQL 5.7 slave server for which 5.7 help content is more appropriate.

This section describes how to manage help table content upgrades when your servers participate in replication. Server versions are one factor in this task. Another is that the help table structure may differ between the master and the slave.

Assume that help content is stored in a file named `fill_help_tables.sql`. In MySQL distributions, this file is located under the `share` or `share/mysql` directory, and the most recent version is always available for download from http://dev.mysql.com/doc/index-other.html.

To upgrade help tables, using the following procedure. Connection parameters are not shown for the `mysql` commands discussed here; in all cases, connect to the server using an account such as `root` that has privileges for modifying tables in the `mysql` database.

1. Upgrade your servers by running `mysql_upgrade`, first on the slaves and then on the master. This is the usual principle of upgrading slaves first.

2. Decide whether you want to replicate help table content from the master to its slaves. If not, load the content on the master and each slave individually. Otherwise, check for and resolve any incompatibilities between help table structure on the master and its slaves, then load the content into the master and let it replicate to the slaves.

   More detail about these two methods of loading help table content follows.

Loading Help Table Content Without Replication to Slaves

To load help table content without replication, run this command on the master and each slave individually, using a `fill_help_tables.sql` file containing content appropriate to the server version:

```
mysql mysql < fill_help_tables.sql
```

Loading Help Table Content With Replication to Slaves

If you do want to replicate help table content, check for help table incompatibilities between your master and its slaves. The `url` column in the `help_category` and `help_topic` tables was originally `CHAR(128)`, but is `TEXT` in newer MySQL versions to accommodate longer URLs. To check help table structure, use this statement:

```
SELECT TABLE_NAME, COLUMN_NAME, COLUMN_TYPE
FROM INFORMATION_SCHEMA.COLUMNS
WHERE TABLE_SCHEMA = 'mysql'
AND COLUMN_NAME = 'url';
```

For tables with the old structure, the statement produces this result:

```
+---------------+-------------+-------------+
| TABLE_NAME    | COLUMN_NAME | COLUMN_TYPE |
+---------------+-------------+-------------+
| help_category | url         | char(128)   |
| help_topic    | url         | char(128)   |
+---------------+-------------+-------------+
```

For tables with the new structure, the statement produces this result:

```
+---------------+-------------+-------------+
| TABLE_NAME    | COLUMN_NAME | COLUMN_TYPE |
+---------------+-------------+-------------+
| help_category | url         | text        |
| help_topic    | url         | text        |
+---------------+-------------+-------------+
```

If the master and slave both have the old structure or both have the new structure, they are compatible and you can replicate help table content by executing this command on the master:

```
mysql mysql < fill_help_tables.sql
```

The table content will load into the master, then replicate to the slaves.

If the master and slave have incompatible help tables (one server has the old structure and the other has the new), you have a choice between not replicating help table content after all, or making the table structures compatible so that you can replicate the content.

- If you decide not to replicate the content after all, upgrade the master and slaves individually using `mysql` with the `--init-command` option, as described previously.
- If instead you decide to make the table structures compatible, upgrade the tables on the server that has the old structure. Suppose that your master server has the old table structure. Upgrade its tables to the new structure manually by executing these statements (binary logging is disabled here to prevent replication of the changes to the slaves, which already have the new structure):

  ```
  SET sql_log_bin=0;
  ALTER TABLE mysql.help_category ALTER COLUMN url TEXT;
  ALTER TABLE mysql.help_topic ALTER COLUMN url TEXT;
  ```

  Then run this command on the master:

  ```
  mysql mysql < fill_help_tables.sql
  ```

  The table content will load into the master, then replicate to the slaves.

18.4.1.30 Replication and Server SQL Mode

Using different server SQL mode settings on the master and the slave may cause the same `INSERT` statements to be handled differently on the master and the slave, leading the master and slave to diverge. For best results, you should always use the same server SQL mode on the master and on the slave. This advice applies whether you are using statement-based or row-based replication.

If you are replicating partitioned tables, using different SQL modes on the master and the slave is likely to cause issues. At a minimum, this is likely to cause the distribution of data among partitions to be

different in the master's and slave's copies of a given table. It may also cause inserts into partitioned tables that succeed on the master to fail on the slave.

For more information, see Section 5.1.8, “Server SQL Modes”. In particular, see SQL Mode Changes in MySQL 5.7, which describes changes in MySQL 5.7 so that you can assess whether your applications will be affected.

18.4.1.31 Replication Retries and Timeouts

The global system variable `slave_transaction_retries` affects replication as follows: If the slave SQL thread fails to execute a transaction because of an `InnoDB` deadlock or because it exceeded the `InnoDB` `innodb_lock_wait_timeout` value, or the `NDB` `TransactionDeadlockDetectionTimeout` or `TransactionInactiveTimeout` value, the slave automatically retries the transaction `slave_transaction_retries` times before stopping with an error. The default value is 10. The total retry count can be seen in the output of `SHOW STATUS`; see Section 5.1.7, “Server Status Variables”.

18.4.1.32 Replication and Time Zones

By default, master and slave servers assume that they are in the same time zone. If you are replicating between servers in different time zones, the time zone must be set on both master and slave. Otherwise, statements depending on the local time on the master are not replicated properly, such as statements that use the `NOW()` or `FROM_UNIXTIME()` functions. Set the time zone in which MySQL server runs by using the `--timezone=`*`timezone_name`* option of the `mysqld_safe` script or by setting the `TZ` environment variable. See also Section 18.4.1.16, “Replication and System Functions”.

18.4.1.33 Replication and Transactions

Mixing transactional and nontransactional statements within the same transaction. In general, you should avoid transactions that update both transactional and nontransactional tables in a replication environment. You should also avoid using any statement that accesses both transactional (or temporary) and nontransactional tables and writes to any of them.

The server uses these rules for binary logging:

- If the initial statements in a transaction are nontransactional, they are written to the binary log immediately. The remaining statements in the transaction are cached and not written to the binary log until the transaction is committed. (If the transaction is rolled back, the cached statements are written to the binary log only if they make nontransactional changes that cannot be rolled back. Otherwise, they are discarded.)
- For statement-based logging, logging of nontransactional statements is affected by the `binlog_direct_non_transactional_updates` system variable. When this variable is `OFF` (the default), logging is as just described. When this variable is `ON`, logging occurs immediately for nontransactional statements occurring anywhere in the transaction (not just initial nontransactional statements). Other statements are kept in the transaction cache and logged when the transaction commits. `binlog_direct_non_transactional_updates` has no effect for row-format or mixed-format binary logging.

Transactional, nontransactional, and mixed statements.
To apply those rules, the server considers a statement nontransactional if it changes only nontransactional tables, and transactional if it changes only transactional tables. In MySQL 8.0, a statement that references both nontransactional and transactional tables and updates *any* of the tables involved, is considered a “mixed” statement. (In previous MySQL release series, a statement that changed both nontransactional and transactional tables was considered mixed.) Mixed statements, like transactional statements, are cached and logged when the transaction commits.

A mixed statement that updates a transactional table is considered unsafe if the statement also performs either of the following actions:

- Updates or reads a transactional table

- Reads a nontransactional table and the transaction isolation level is less than REPEATABLE_READ

A mixed statement following the update of a transactional table within a transaction is considered unsafe if it performs either of the following actions:

- Updates any table and reads from any temporary table
- Updates a nontransactional table and binlog_direct_non_trans_update is OFF

For more information, see Section 18.2.1.3, “Determination of Safe and Unsafe Statements in Binary Logging”.

Note

A mixed statement is unrelated to mixed binary logging format.

In situations where transactions mix updates to transactional and nontransactional tables, the order of statements in the binary log is correct, and all needed statements are written to the binary log even in case of a `ROLLBACK`. However, when a second connection updates the nontransactional table before the first connection transaction is complete, statements can be logged out of order because the second connection update is written immediately after it is performed, regardless of the state of the transaction being performed by the first connection.

Using different storage engines on master and slave. It is possible to replicate transactional tables on the master using nontransactional tables on the slave. For example, you can replicate an `InnoDB` master table as a `MyISAM` slave table. However, if you do this, there are problems if the slave is stopped in the middle of a `BEGIN` ... `COMMIT` block because the slave restarts at the beginning of the `BEGIN` block.

In MySQL 8.0, it is also safe to replicate transactions from `MyISAM` tables on the master to transactional tables—such as tables that use the `InnoDB` storage engine—on the slave. In such cases, an `AUTOCOMMIT=1` statement issued on the master is replicated, thus enforcing `AUTOCOMMIT` mode on the slave.

When the storage engine type of the slave is nontransactional, transactions on the master that mix updates of transactional and nontransactional tables should be avoided because they can cause inconsistency of the data between the master transactional table and the slave nontransactional table. That is, such transactions can lead to master storage engine-specific behavior with the possible effect of replication going out of synchrony. MySQL does not issue a warning about this currently, so extra care should be taken when replicating transactional tables from the master to nontransactional tables on the slaves.

Changing the binary logging format within transactions. The `binlog_format` system variable is read-only as long as a transaction is in progress.

Every transaction (including `autocommit` transactions) is recorded in the binary log as though it starts with a `BEGIN` statement, and ends with either a `COMMIT` or a `ROLLBACK` statement. In MySQL 8.0, this true is even for statements affecting tables that use a nontransactional storage engine (such as `MyISAM`).

18.4.1.34 Replication and Transaction Inconsistencies

Inconsistencies in the sequence of transactions that have been executed from the relay log can occur depending on your replication configuration. This section explains how to avoid inconsistencies and solve any problems they cause.

The following types of inconsistencies can exist:

- *Half-applied transactions*. A transaction which updates non-transactional tables has applied some but not all of its changes.

- *Gaps*. A gap is a transaction that has not been (fully) applied, even though some later transaction has been applied. Gaps can only appear when using a multi-threaded slave. To avoid gaps occurring, set `slave_preserve_commit_order=1`, which requires `slave_parallel_type=LOGICAL_CLOCK`, and that `log-bin` and `log-slave-updates` are also enabled.

- *Gap-free low-watermark position*. Even in the absence of gaps, it is possible that transactions after `Exec_master_log_pos` have not been applied. That is, all transactions up to point `N` have been applied, and no transactions after `N` have been applied, but `Exec_master_log_pos` has a value smaller than `N.` This can only happen on multi-threaded slaves. Enabling `slave_preserve_commit_order` does *not* prevent gap-free low-watermark positions.

The following scenarios are relevant to the existence of half-applied transactions, gaps, and gap-free low-watermark position inconsistencies:

1. While slave threads are running, there may be gaps and half-applied transactions.

2. `mysqld` shuts down. Both clean and unclean shutdown abort ongoing transactions and may leave gaps and half-applied transactions.

3. `KILL` of replication threads (the SQL thread when using a single-threaded slave, the coordinator thread when using a multi-threaded slave). This aborts ongoing transactions and may leave gaps and half-applied transactions.

4. Error in applier threads. This may leave gaps. If the error is in a mixed transaction, that transaction is half-applied. When using a multi-threaded slave, workers which have not received an error complete their queues, so it may take time to stop all threads.

5. `STOP SLAVE` when using a multi-threaded slave. After issuing `STOP SLAVE`, the slave waits for any gaps to be filled and then updates `Exec_master_log_pos`. This ensures it never leaves gaps or gap-free low-watermark positions, unless any of the cases above applies (in other words, before `STOP SLAVE` completes, either an error happens, or another thread issues `KILL`, or the server restarts. In these cases, `STOP SLAVE` returns successfully.)

6. If the last transaction in the relay log is only half-received and the multi-threaded slave coordinator has started to schedule the transaction to a worker, then `STOP SLAVE` waits up to 60 seconds for the transaction to be received. After this timeout, the coordinator gives up and aborts the transaction. If the transaction is mixed, it may be left half-completed.

7. `STOP SLAVE` when using a single-threaded slave. If the ongoing transaction only updates transactional tables, it is rolled back and `STOP SLAVE` stops immediately. If the ongoing transaction is mixed, `STOP SLAVE` waits up to 60 seconds for the transaction to complete. After this timeout, it aborts the transaction, so it may be left half-completed.

The global variable `rpl_stop_slave_timeout` is unrelated to the process of stopping the replication threads. It only makes the client that issues `STOP SLAVE` return to the client, but the replication threads continue to try to stop.

If a replication channel has gaps, it has the following consequences:

1. The slave database is in a state that may never have existed on the master.

2. The field `Exec_master_log_pos` in `SHOW SLAVE STATUS` is only a "low-watermark". In other words, transactions appearing before the position are guaranteed to have committed, but transactions after the position may have committed or not.

3. `CHANGE MASTER TO` statements for that channel fail with an error, unless the applier threads are running and the `CHANGE MASTER TO` statement only sets receiver options.

4. If `mysqld` is started with `--relay-log-recovery`, no recovery is done for that channel, and a warning is printed.

5. If `mysqldump` is used with `--dump-slave`, it does not record the existence of gaps; thus it prints `CHANGE MASTER TO` with `RELAY_LOG_POS` set to the low-watermark position in `Exec_master_log_pos`.

   After applying the dump on another server, and starting the replication threads, transactions appearing after the position are replicated again. Note that this is harmless if GTIDs are enabled (however, in that case it is not recommended to use `--dump-slave`).

If a replication channel has a gap-free low-watermark position, cases 2 to 5 above apply, but case 1 does not.

The gap-free low-watermark position information is persisted in binary format in the internal table `mysql.slave_worker_info`. `START SLAVE [SQL_THREAD]` always consults this information so that it applies only the correct transactions. This remains true even if `slave_parallel_workers` has been changed to 0 before `START SLAVE`, and even if `START SLAVE` is used with `UNTIL` clauses. `START SLAVE UNTIL SQL_AFTER_MTS_GAPS` only applies as many transactions as needed in order to fill in the gaps. If `START SLAVE` is used with `UNTIL` clauses that tell it to stop before it has consumed all the gaps, then it leaves remaining gaps.

Warning

`RESET SLAVE` removes the relay logs and resets the replication position. Thus issuing `RESET SLAVE` on a slave with gaps means the slave loses any information about the gaps, without correcting the gaps.

`slave-preserve-commit-order` ensures that there are no gaps. However, it is still possible that `Exec_master_log_pos` is just a gap-free low-watermark position in scenarios 1 to 4 above. That is, there may be transactions after `Exec_master_log_pos` which have been applied. Therefore the cases numbered 2 to 5 above (but not case 1) apply, even when `slave-preserve-commit-order` is enabled.

18.4.1.35 Replication and Triggers

With statement-based replication, triggers executed on the master also execute on the slave. With row-based replication, triggers executed on the master do not execute on the slave. Instead, the row changes on the master resulting from trigger execution are replicated and applied on the slave.

This behavior is by design. If under row-based replication the slave applied the triggers as well as the row changes caused by them, the changes would in effect be applied twice on the slave, leading to different data on the master and the slave.

If you want triggers to execute on both the master and the slave—perhaps because you have different triggers on the master and slave—you must use statement-based replication. However, to enable slave-side triggers, it is not necessary to use statement-based replication exclusively. It is sufficient to switch to statement-based replication only for those statements where you want this effect, and to use row-based replication the rest of the time.

A statement invoking a trigger (or function) that causes an update to an `AUTO_INCREMENT` column is not replicated correctly using statement-based replication. MySQL 8.0 marks such statements as unsafe. (Bug #45677)

A trigger can have triggers for different combinations of trigger event (`INSERT`, `UPDATE`, `DELETE`) and action time (`BEFORE`, `AFTER`), and multiple triggers are permitted.

For brevity, “multiple triggers” here is shorthand for “multiple triggers that have the same trigger event and action time.”

Upgrades. Suppose that you upgrade an old server that does not support multiple triggers (lower than 5.7.2) to a version that does (MySQL 5.7.2 or higher). If the new server is a replication master and has old slaves that do not support multiple triggers, an error occurs on those slaves if a trigger is created on

the master for a table that already has a trigger with the same trigger event and action time. To avoid this problem, upgrade the slaves first, then upgrade the master.

Downgrades. If you downgrade a server that supports multiple triggers to an older version that does not, the downgrade has these effects:

- For each table that has triggers, all trigger definitions are in the `.TRG` file for the table. However, if there are multiple triggers with the same trigger event and action time, the server executes only one of them when the trigger event occurs. For information about `.TRG` files, see Table Trigger Storage.
- If triggers for the table are added or dropped subsequent to the downgrade, the server rewrites the table's `.TRG` file. The rewritten file retains only one trigger per combination of trigger event and action time; the others are lost.

To avoid these problems, modify your triggers before downgrading. For each table that has multiple triggers per combination of trigger event and action time, convert each such set of triggers to a single trigger as follows:

1. For each trigger, create a stored routine that contains all the code in the trigger. Values accessed using `NEW` and `OLD` can be passed to the routine using parameters. If the trigger needs a single result value from the code, you can put the code in a stored function and have the function return the value. If the trigger needs multiple result values from the code, you can put the code in a stored procedure and return the values using `OUT` parameters.
2. Drop all triggers for the table.
3. Create one new trigger for the table that invokes the stored routines just created. The effect for this trigger is thus the same as the multiple triggers it replaces.

18.4.1.36 Replication and TRUNCATE TABLE

`TRUNCATE TABLE` is normally regarded as a DML statement, and so would be expected to be logged and replicated using row-based format when the binary logging mode is `ROW` or `MIXED`. However this caused issues when logging or replicating, in `STATEMENT` or `MIXED` mode, tables that used transactional storage engines such as `InnoDB` when the transaction isolation level was `READ COMMITTED` or `READ UNCOMMITTED`, which precludes statement-based logging.

`TRUNCATE TABLE` is treated for purposes of logging and replication as DDL rather than DML so that it can be logged and replicated as a statement. However, the effects of the statement as applicable to `InnoDB` and other transactional tables on replication slaves still follow the rules described in Section 13.1.30, “TRUNCATE TABLE Syntax” governing such tables. (Bug #36763)

18.4.1.37 Replication and User Name Length

The maximum length of MySQL user names is 32 characters in MySQL 5.7.8 or higher. Replication of user names longer than 16 characters to a slave that supports only shorter user names will fail. However, this should occur only when replicating from a newer master to an older slave, which is not a recommended configuration.

18.4.1.38 Replication and Variables

System variables are not replicated correctly when using `STATEMENT` mode, except for the following variables when they are used with session scope:

- `auto_increment_increment`
- `auto_increment_offset`
- `character_set_client`
- `character_set_connection`

- character_set_database
- character_set_server
- collation_connection
- collation_database
- collation_server
- foreign_key_checks
- identity
- last_insert_id
- lc_time_names
- pseudo_thread_id
- sql_auto_is_null
- time_zone
- timestamp
- unique_checks

When MIXED mode is used, the variables in the preceding list, when used with session scope, cause a switch from statement-based to row-based logging. See Section 5.4.4.3, “Mixed Binary Logging Format”.

sql_mode is also replicated except for the NO_DIR_IN_CREATE mode; the slave always preserves its own value for NO_DIR_IN_CREATE, regardless of changes to it on the master. This is true for all replication formats.

However, when mysqlbinlog parses a SET @@sql_mode = *mode* statement, the full *mode* value, including NO_DIR_IN_CREATE, is passed to the receiving server. For this reason, replication of such a statement may not be safe when STATEMENT mode is in use.

The default_storage_engine system variable is not replicated, regardless of the logging mode; this is intended to facilitate replication between different storage engines.

The read_only system variable is not replicated. In addition, the enabling this variable has different effects with regard to temporary tables, table locking, and the SET PASSWORD statement in different MySQL versions.

The max_heap_table_size system variable is not replicated. Increasing the value of this variable on the master without doing so on the slave can lead eventually to Table is full errors on the slave when trying to execute INSERT statements on a MEMORY table on the master that is thus permitted to grow larger than its counterpart on the slave. For more information, see Section 18.4.1.23, “Replication and MEMORY Tables”.

In statement-based replication, session variables are not replicated properly when used in statements that update tables. For example, the following sequence of statements will not insert the same data on the master and the slave:

```
SET max_join_size=1000;
INSERT INTO mytable VALUES(@@max_join_size);
```

This does not apply to the common sequence:

```
SET time_zone=...;
INSERT INTO mytable VALUES(CONVERT_TZ(..., ..., @@time_zone));
```

Replication of session variables is not a problem when row-based replication is being used, in which case, session variables are always replicated safely. See Section 18.2.1, “Replication Formats”.

In MySQL 8.0, the following session variables are written to the binary log and honored by the replication slave when parsing the binary log, regardless of the logging format:

- `sql_mode`
- `foreign_key_checks`
- `unique_checks`
- `character_set_client`
- `collation_connection`
- `collation_database`
- `collation_server`
- `sql_auto_is_null`

Important

Even though session variables relating to character sets and collations are written to the binary log, replication between different character sets is not supported.

To help reduce possible confusion, we recommend that you always use the same setting for the `lower_case_table_names` system variable on both master and slave, especially when you are running MySQL on platforms with case-sensitive file systems.

18.4.1.39 Replication and Views

Views are always replicated to slaves. Views are filtered by their own name, not by the tables they refer to. This means that a view can be replicated to the slave even if the view contains a table that would normally be filtered out by `replication-ignore-table` rules. Care should therefore be taken to ensure that views do not replicate table data that would normally be filtered for security reasons.

Replication from a table to a same-named view is supported using statement-based logging, but not when using row-based logging. Trying to do so when row-based logging is in effect causes an error.

18.4.2 Replication Compatibility Between MySQL Versions

MySQL supports replication from one release series to the next higher release series. For example, you can replicate from a master running MySQL 5.6 to a slave running MySQL 5.7, from a master running MySQL 5.7 to a slave running MySQL 8.0, and so on.

However, you may encounter difficulties when replicating from an older master to a newer slave if the master uses statements or relies on behavior no longer supported in the version of MySQL used on the slave. For example, in MySQL 5.5, `CREATE TABLE ... SELECT` statements are permitted to change tables other than the one being created, but are no longer allowed to do so in MySQL 5.6 (see Section 18.4.1.6, “Replication of CREATE TABLE ... SELECT Statements”).

The use of more than two MySQL Server versions is not supported in replication setups involving multiple masters, regardless of the number of master or slave MySQL servers. This restriction applies not only to release series, but to version numbers within the same release series as well. For example, if you are using a chained or circular replication setup, you cannot use MySQL 8.0.1, MySQL 8.0.2, and MySQL 8.0.4 concurrently, although you could use any two of these releases together.

Important

It is strongly recommended to use the most recent release available within a given MySQL release series because replication (and other) capabilities are continually being improved. It is also recommended to upgrade masters and slaves that use early releases of a release series of MySQL to GA (production) releases when the latter become available for that release series.

Replication from newer masters to older slaves may be possible, but is generally not supported. This is due to a number of factors:

- **Binary log format changes.** The binary log format can change between major releases. While we attempt to maintain backward compatibility, this is not always possible.

  This also has significant implications for upgrading replication servers; see Section 18.4.3, “Upgrading a Replication Setup”, for more information.

- For more information about row-based replication, see Section 18.2.1, “Replication Formats”.

- **SQL incompatibilities.** You cannot replicate from a newer master to an older slave using statement-based replication if the statements to be replicated use SQL features available on the master but not on the slave.

  However, if both the master and the slave support row-based replication, and there are no data definition statements to be replicated that depend on SQL features found on the master but not on the slave, you can use row-based replication to replicate the effects of data modification statements even if the DDL run on the master is not supported on the slave.

For more information on potential replication issues, see Section 18.4.1, “Replication Features and Issues”.

18.4.3 Upgrading a Replication Setup

When you upgrade servers that participate in a replication setup, the procedure for upgrading depends on the current server versions and the version to which you are upgrading. This section provides information about how upgrading affects replication. For general information about upgrading MySQL, see Section 2.10.1, “Upgrading MySQL”

When you upgrade a master to 8.0 from an earlier MySQL release series, you should first ensure that all the slaves of this master are using the same 8.0.x release. If this is not the case, you should first upgrade the slaves. To upgrade each slave, shut it down, upgrade it to the appropriate 8.0.x version, restart it, and restart replication. Relay logs created by the slave after the upgrade are in 8.0 format.

Changes affecting operations in strict SQL mode may result in replication failure on an updated slave. For example, the server restricts insertion of a `DEFAULT` value of 0 for temporal data types in strict mode (`STRICT_TRANS_TABLES` or `STRICT_ALL_TABLES`). A resulting incompatibility for replication if you use statement-based logging (`binlog_format=STATEMENT`) is that if a slave is upgraded, a nonupgraded master will execute statements without error that may fail on the slave and replication will stop. To deal with this, stop all new statements on the master and wait until the slaves catch up. Then upgrade the slaves. Alternatively, if you cannot stop new statements, temporarily change to row-based logging on the master (`binlog_format=ROW`) and wait until all slaves have processed all binary logs produced up to the point of this change. Then upgrade the slaves.

After the slaves have been upgraded, shut down the master, upgrade it to the same 8.0.x release as the slaves, and restart it. If you had temporarily changed the master to row-based logging, change it back to statement-based logging. The 8.0 master is able to read the old binary logs written prior to the upgrade and to send them to the 8.0 slaves. The slaves recognize the old format and handle it properly. Binary logs created by the master subsequent to the upgrade are in 8.0 format. These too are recognized by the 8.0 slaves.

In other words, when upgrading to MySQL 8.0, the slaves must be MySQL 8.0 before you can upgrade the master to 8.0. Note that downgrading from 8.0 to older versions does not work so simply: You must ensure that any 8.0 binary log or relay log has been fully processed, so that you can remove it before proceeding with the downgrade.

Some upgrades may require that you drop and re-create database objects when you move from one MySQL series to the next. For example, collation changes might require that table indexes be rebuilt. Such operations, if necessary, are detailed at Section 2.10.1.1, “Changes Affecting Upgrades to MySQL 8.0”. It is safest to perform these operations separately on the slaves and the master, and to disable replication of these operations from the master to the slave. To achieve this, use the following procedure:

1. Stop all the slaves and upgrade them. Restart them with the `--skip-slave-start` option so that they do not connect to the master. Perform any table repair or rebuilding operations needed to re-create database objects, such as use of `REPAIR TABLE` or `ALTER TABLE`, or dumping and reloading tables or triggers.
2. Disable the binary log on the master. To do this without restarting the master, execute a `SET sql_log_bin = 0` statement. Alternatively, stop the master and restart it without the `--log-bin` option. If you restart the master, you might also want to disallow client connections. For example, if all clients connect using TCP/IP, use the `--skip-networking` option when you restart the master.
3. With the binary log disabled, perform any table repair or rebuilding operations needed to re-create database objects. The binary log must be disabled during this step to prevent these operations from being logged and sent to the slaves later.
4. Re-enable the binary log on the master. If you set `sql_log_bin` to 0 earlier, execute a `SET sql_log_bin = 1` statement. If you restarted the master to disable the binary log, restart it with `--log-bin`, and without `--skip-networking` so that clients and slaves can connect.
5. Restart the slaves, this time without the `--skip-slave-start` option.

If you are upgrading an existing replication setup from a version of MySQL that does not support global transaction identifiers to a version that does, you should not enable GTIDs on either the master or the slave before making sure that the setup meets all the requirements for GTID-based replication. For example `server_uuid`, which was added in MySQL 5.6, must exist for GTIDs to function correctly. See Section 18.1.3.2, “Setting Up Replication Using GTIDs”, which contains information about converting existing replication setups to use GTID-based replication.

18.4.4 Troubleshooting Replication

If you have followed the instructions but your replication setup is not working, the first thing to do is *check the error log for messages*. Many users have lost time by not doing this soon enough after encountering problems.

If you cannot tell from the error log what the problem was, try the following techniques:

- Verify that the master has binary logging enabled by issuing a `SHOW MASTER STATUS` statement. If logging is enabled, `Position` is nonzero. If binary logging is not enabled, verify that you are running the master with the `--log-bin` option.
- Verify that the master and slave both were started with the `--server-id` option and that the ID value is unique on each server.
- Verify that the slave is running. Use `SHOW SLAVE STATUS` to check whether the `Slave_IO_Running` and `Slave_SQL_Running` values are both `Yes`. If not, verify the options that were used when starting the slave server. For example, `--skip-slave-start` prevents the slave threads from starting until you issue a `START SLAVE` statement.
- If the slave is running, check whether it established a connection to the master. Use `SHOW PROCESSLIST`, find the I/O and SQL threads and check their `State` column to see what they

display. See Section 18.2.2, “Replication Implementation Details”. If the I/O thread state says `Connecting to master`, check the following:

- Verify the privileges for the user being used for replication on the master.
- Check that the host name of the master is correct and that you are using the correct port to connect to the master. The port used for replication is the same as used for client network communication (the default is `3306`). For the host name, ensure that the name resolves to the correct IP address.
- Check that networking has not been disabled on the master or slave. Look for the `skip-networking` option in the configuration file. If present, comment it out or remove it.
- If the master has a firewall or IP filtering configuration, ensure that the network port being used for MySQL is not being filtered.
- Check that you can reach the master by using `ping` or `traceroute`/`tracert` to reach the host.

- If the slave was running previously but has stopped, the reason usually is that some statement that succeeded on the master failed on the slave. This should never happen if you have taken a proper snapshot of the master, and never modified the data on the slave outside of the slave thread. If the slave stops unexpectedly, it is a bug or you have encountered one of the known replication limitations described in Section 18.4.1, “Replication Features and Issues”. If it is a bug, see Section 18.4.5, “How to Report Replication Bugs or Problems”, for instructions on how to report it.

- If a statement that succeeded on the master refuses to run on the slave, try the following procedure if it is not feasible to do a full database resynchronization by deleting the slave's databases and copying a new snapshot from the master:

  1. Determine whether the affected table on the slave is different from the master table. Try to understand how this happened. Then make the slave's table identical to the master's and run `START SLAVE`.
  2. If the preceding step does not work or does not apply, try to understand whether it would be safe to make the update manually (if needed) and then ignore the next statement from the master.
  3. If you decide that the slave can skip the next statement from the master, issue the following statements:

     ```
     mysql> SET GLOBAL sql_slave_skip_counter = N;
     mysql> START SLAVE;
     ```

     The value of *`N`* should be 1 if the next statement from the master does not use `AUTO_INCREMENT` or `LAST_INSERT_ID()`. Otherwise, the value should be 2. The reason for using a value of 2 for statements that use `AUTO_INCREMENT` or `LAST_INSERT_ID()` is that they take two events in the binary log of the master.

     See also Section 13.4.2.5, “SET GLOBAL sql_slave_skip_counter Syntax”.
  4. If you are sure that the slave started out perfectly synchronized with the master, and that no one has updated the tables involved outside of the slave thread, then presumably the discrepancy is the result of a bug. If you are running the most recent version of MySQL, please report the problem. If you are running an older version, try upgrading to the latest production release to determine whether the problem persists.

18.4.5 How to Report Replication Bugs or Problems

When you have determined that there is no user error involved, and replication still either does not work at all or is unstable, it is time to send us a bug report. We need to obtain as much information as

possible from you to be able to track down the bug. Please spend some time and effort in preparing a good bug report.

If you have a repeatable test case that demonstrates the bug, please enter it into our bugs database using the instructions given in Section 1.7, “How to Report Bugs or Problems”. If you have a “phantom” problem (one that you cannot duplicate at will), use the following procedure:

1. Verify that no user error is involved. For example, if you update the slave outside of the slave thread, the data goes out of synchrony, and you can have unique key violations on updates. In this case, the slave thread stops and waits for you to clean up the tables manually to bring them into synchrony. *This is not a replication problem. It is a problem of outside interference causing replication to fail.*

2. Run the slave with the `--log-slave-updates` and `--log-bin` options. These options cause the slave to log the updates that it receives from the master into its own binary logs.

3. Save all evidence before resetting the replication state. If we have no information or only sketchy information, it becomes difficult or impossible for us to track down the problem. The evidence you should collect is:

   - All binary log files from the master
   - All binary log files from the slave
   - The output of `SHOW MASTER STATUS` from the master at the time you discovered the problem
   - The output of `SHOW SLAVE STATUS` from the slave at the time you discovered the problem
   - Error logs from the master and the slave

4. Use `mysqlbinlog` to examine the binary logs. The following should be helpful to find the problem statement. *`log_file`* and *`log_pos`* are the `Master_Log_File` and `Read_Master_Log_Pos` values from `SHOW SLAVE STATUS`.

   ```
   shell> mysqlbinlog --start-position=log_pos log_file | head
   ```

After you have collected the evidence for the problem, try to isolate it as a separate test case first. Then enter the problem with as much information as possible into our bugs database using the instructions at Section 1.7, “How to Report Bugs or Problems”.

Chapter 19 MySQL Shell User Guide

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This section provides user documentation for MySQL Shell, which enables you to prototype code using the X DevAPI to communicate with a MySQL Server running the X Plugin.

MySQL Shell is an advanced command-line client and code editor for the MySQL Server. In addition to SQL, MySQL Shell also offers scripting capabilities for JavaScript and Python. When MySQL Shell is connected to the MySQL Server through the X Protocol, the X DevAPI can be used to work with both relational and document data.

19.1 MySQL Shell Features

The following features are available in MySQL Shell.

Interactive Code Execution

MySQL Shell provides an interactive code execution mode, where you type code at the MySQL Shell prompt and each entered statement is processed, with the result of the processing printed onscreen.

Supported Languages

MySQL Shell processes code in the following languages: JavaScript, Python and SQL. Any entered code is processed as one of these languages, based on the language that is currently active. There are also specific commands that are treated as “shell commands” which enable you to configure MySQL Shell or retrieve information. For more information see Section 19.4.1, “MySQL Shell Commands”.

Batch Code Execution

In addition to the interactive execution of code, MySQL Shell can also take code from different sources and process it. This method of processing code in a non-interactive way is called *Batch Execution*.

As batch execution mode is intended for script processing of a single language, it is limited to having minimal non-formatted output and disabling the execution of commands. To avoid these limitations, use

the `--interactive` command-line option, which tells MySQL Shell to execute the input as if it were an interactive session. In this mode the input is processed *line by line* just as if each line were typed in an interactive session. For more information see Section 19.3.5, “Batch Mode Made Interactive”.

Output Formats

MySQL Shell provides output in different formats depending on how it is used: Tabbed, Table and JSON. For more information see Section 19.3.3, “Output Formats”.

Multiple-line Support

Multiple-line code can be written using a command, enabling MySQL Shell to cache multiple lines and then execute them as a single statement. For more information see Section 19.3.5.1, “Multiple-line Support”.

Stored Sessions

MySQL Shell maintains a list of connection data for different servers. You can connect to a server using these stored sessions. The list of stored sessions can be edited. See Section 19.5, “Stored Sessions” for more information.

Application Log

MySQL Shell can be configured to log information about the execution process. For more information see Section 19.6, “MySQL Shell Application Log”.

X DevAPI

Code written in MySQL Shell uses the X DevAPI, provided for Python and JavaScript. For general information, see X DevAPI User Guide.

X Protocol Support

MySQL Shell is designed to provide an integrated command-line client for all MySQL products which support X Protocol. The development features of MySQL Shell are designed for sessions using the X Protocol. MySQL Shell can also connect to MySQL Servers that do not support the X Protocol using the legacy MySQL Protocol. A minimal set of features from the X DevAPI are available for sessions created using the legacy MySQL protocol.

Global Session

Interaction with a MySQL Server is done through a Session object. For Python and JavaScript, a Session can be created through the `getSession` and `getNodeSession` functions of the mysqlx module. If a session is created in JavaScript mode using any of these methods, it is available only in JavaScript mode. The same happens if the session is created in Python mode. None of these sessions can be used in SQL mode.

For SQL Mode, the concept of Global Session is supported by the MySQL Shell. A Global Session is created when the connection information is passed to MySQL Shell using command-line options, or by using the `\connect` command.

The Global Session is used to execute statements in SQL mode and the same session is available in both Python or JavaScript modes. When a Global Session is created, a variable called `session` is set in the scripting languages, so you can execute code in the different languages by switching the active mode.

For more information, see Section 19.2.1, “MySQL Shell Sessions”.

19.2 Getting Started with MySQL Shell

This section describes how to get started with MySQL Shell. This section assumes you have a MySQL Server running X Plugin and that you have installed MySQL Shell, see Section 20.3, “Setting Up MySQL as a Document Store”.

19.2.1 MySQL Shell Sessions

This section explains the different types of sessions in MySQL Shell and how to create and configure them.

19.2.1.1 MySQL Shell Sessions Explained

MySQL Shell is a unified interface to operate MySQL Server through scripting languages such as JavaScript or Python. To maintain compatibility with previous versions, SQL can also be executed in certain modes. A connection to a MySQL server is required. In MySQL Shell these connections are handled by a *Session* object.

The following different types of Session objects are available:

- *XSession*: Use this session type for new application development. It offers the best integration with MySQL Server, and therefore, it is used by default. SQL execution is not supported and therefore it is not compatible with MySQL Shell's *SQL Mode*.
- *Node Session*: Use this session type for SQL execution on a MySQL Server with the X Protocol enabled. SQL execution is available with this session type, therefore it can be used in MySQL Shell's *SQL Mode*.

  This session type should only be used when connecting *directly* to an X Protocol enabled MySQL Server.
- *Classic Session* Use this session type to interact with MySQL Servers that do not have the X Protocol enabled. SQL execution is available with this session type, therefore it can be used in MySQL Shell's *SQL Mode*.

  The development API available for this type of session is very limited. For example, there are no CRUD operations, no collection handling, and binding is not supported.

19.2.1.2 Choosing a MySQL Shell Session Type

MySQL Shell creates an XSession object by default. To choose which type of session should be created, use one of these options:

- `--node` creates a Node Session.
- `--classic` creates a Classic Session.
- `--x` creates an XSession.

For more information, see Section 19.2.2.1, “Connecting using a URI String” and Section 19.2.2.2, “Connecting using Individual Parameters”.

19.2.2 MySQL Shell Connections

MySQL Shell can be configured to connect to a MySQL Server running X Plugin using command-line options at start up or from MySQL Shell itself. The address of the MySQL Server can be specified using individual parameters, such as user, hostname and port, or using a Uniform Resource Identifier (URI) in the format `user@host:port/schema`, such as `mike@myserver:33060/testDB`. The following sections describe these methods.

Regardless of the method you choose to connect it is important to understand how passwords are handled by MySQL Shell:

Important

By default connections are assumed to require a password. The password is requested at the login prompt. To specify a password-less account use the `--password` option, or use a `:` after the `user` in a URI.

If you do not specify parameters for a connection the following defaults are used:

- user defaults to the current system user name
- host defaults to localhost
- port defaults to the X Plugin port 33060 when using an X Session, and port 3306 when using a Classic session

MySQL Shell connections using X Protocol *always* use TCP, using Linux sockets is not supported. MySQL Shell connections using MySQL Protocol can use Linux sockets by specifying the `--socket` command-line option.

19.2.2.1 Connecting using a URI String

You configure the MySQL Server which MySQL Shell connects to by passing the connection data in string format using the `--uri` command-line option.

Use the following format:

```
[dbuser[:[dbpassword]]@]host[:port][/schema]
```

Descriptions of these options:

- *`dbuser`*: specifies the MySQL user account to be used for the authentication process.
- *`dbpassword`*: specifies the dbuser password to be used for the authentication process.

Warning

Storing the password in the URI is not recommended.

- *`host`*: specifies the host to which the session object connects. If not specified, *localhost* is used by default.
- *`port`*: specifies which port the target MySQL server is listening on for the connection. If not specified, 33060 is used by default for the X Protocol enabled sessions, and 3306 is the default for traditional MySQL protocol sessions.
- *`schema`*: Indicates the database to be set as default when the session is established.

If no password is specified using the URI, which is recommended, then the password is prompted for. The following examples show how to use these command-line options:

- Connect with a Node Session at port 33065.

```
shell> mysqlsh --uri user@localhost:33065 --node
```

- Connect with a Classic Session.

```
shell> mysqlsh --uri user@localhost --classic
```

Although using a password-less account is not recommended, you can specify a user without a password using a `:` after the user name, for example:

```
shell> mysqlsh --uri user:@localhost
```

In addition to specifying a URI manually, you can use a stored session. See Section 19.5, “Stored Sessions” for more information.

19.2.2.2 Connecting using Individual Parameters

In addition to specifying connection parameters using a URI, it is also possible to define the connection data using separate parameters for each value.

Use the following parameters:

- `--dbuser` (`-u`) *`value`*
- `--dbpassword` *`value`*
- `--host` (`-h`) *`value`*
- `--port` (`-P`) *`value`*
- `--schema` (`-D`) *`value`*
- `--session-type` *`value`*. Removed in version 1.0.4.
- `--password` (`-p`)

The first 5 parameters match the tokens used in the URI format described at Section 19.2.2.1, “Connecting using a URI String”.

The `--password` parameter indicates the user should connect *without* a password.

For consistency, the following aliases are supported for some parameters:

- `--user` is equivalent to `--dbuser`
- `--password` is equivalent to `--dbpassword`
- `--database` is equivalent to `--schema`

Attempt to establish an XSession with a specified user at port 33065.

```
shell> mysqlsh -u user -h localhost -P 33065
```

Attempt to establish a Classic Session with a specified user.

```
shell> mysqlsh -u user -h localhost --classic
```

Attempt to establish a Node Session with a specified user.

```
shell> mysqlsh --node -u user -h localhost
```

19.2.2.3 Using SSL for Secure Connections

Using SSL is possible when connecting to an SSL enabled MySQL server.

To configure an SSL connection, use the following command-line options:

- `--ssl` : This enables or disables connections through SSL. If set to 0, the other SSL command-line options are ignored.
- `--ssl-ca=filename`: The path to a file in PEM format that contains a list of trusted SSL certificate authorities.
- `--ssl-cert=filename`: The name of the SSL certificate file in PEM format to use for establishing a secure connection.
- `--ssl-key=filename`: The name of the SSL key file in PEM format to use for establishing a secure connection.

The `--ssl` option is assumed to be 1 (enabled) if the other SSL options are set.

19.2.2.4 Connecting using both URI and Individual Parameters

When the `--uri` option is specified in combination with some of the individual parameters, the address specified by the `--uri` option is used as the base connection data and the values provided using individual parameters override the corresponding value from the URI. If the `--user` option is used, it would replace any user specified as part of a URI.

For example to establish an XSession and override *`user`* from the URI:

```
shell> mysqlsh --uri user@localhost:33065 --user otheruser
```

19.2.2.5 Creating a Session Using Shell Commands

If you open MySQL Shell without specifying connection parameters, MySQL Shell opens without an established global session. It is possible to establish a global session once MySQL Shell has been started using the following *Shell Commands*:

- `\connect URI`: Creates an XSession.
- `\connect -n URI`: Creates a Node Session.
- `\connect -c URI`: Creates a Classic Session.

Configure the connection using the URI parameter, which follows the same syntax as for the `--uri` command-line option. For additional information, see Section 19.2.2.1, “Connecting using a URI String”.

For example:

```
mysql-js> \connect root@localhost
Creating XSession to root@localhost...
Enter password: ****
No default schema selected.
mysql-js>
```

Creating a Secure Session Using SSL

To establish an SSL connection, the URI parameter passed to the connect shell commands must include the SSL information as URL parameters. For example:

```
mysql-js> \connect root@localhost?ssl_ca=/path/to/ca/file&\
ssl_cert=/path/to/cert/file&ssl_key=/path/to/key/file
Creating XSession to root@localhost...
Enter password: ****
```

```
No default schema selected.
mysql-js>
```

19.2.2.6 Connections in JavaScript and Python

When a connection is made using the command-line options or by using any of the shell commands, a global session object is created. This session is global because once created, it can be used in any of the MySQL Shell execution modes.

Any global session object is available in JavaScript or Python modes because a variable called **session** holds a reference to it.

In SQL mode, both Node and Classic sessions can be used because they both expose SQL execution. In SQL mode an XSession cannot be used.

In addition to the global session object, sessions can be established and assigned to a different variable by using the functions available in the mysql and mysqlx JavaScript and Python modules.

For example, the following functions are provided by these modules:

- **mysql.getSession(connectionData[, password])**
- **mysqlx.getNodeSession(connectionData[, password])**
- **mysql.getClassicSession(connectionData[, password])**

The first of these functions is used to create an XSession which features the most comprehensive development API and supports X Protocol.

The second creates a Node Session which connects to a X Protocol enabled MySQL Server and allows SQL Execution.

The latter returns a Classic Session object which uses the traditional MySQL protocol and has a very limited development API.

connectionData can be either a URI as specified above or a dictionary containing the connection parameters. See Section 19.2.2.1, “Connecting using a URI String”.

The following example shows how to create a Node Session using the X Protocol:

```
mysql-js> var mysqlx=require('mysqlx').mysqlx;

mysql-js> var session=mysqlx.getNodeSession('root@localhost');

mysql-js> print(session)
<NodeSession:root@localhost>
mysql-js>
```

The following example shows how to create a Node Session using the X Protocol so that you can execute SQL:

```
mysql-js> var mysqlx=require('mysqlx').mysqlx;

mysql-js> var session=mysqlx.getNodeSession({host: 'localhost', dbUser: 'root'});

mysql-js> print(session)
<NodeSession:root@localhost>
mysql-js>
```

The following example shows how to create a Classic Session:

```
```

```
mysql-js> var mysql=require('mysql').mysql;

mysql-js> var session = mysql.getClassicSession('root@localhost:3307');
mysql-js> print(session)
<ClassicSession:root@localhost:3307>
mysql-js>
```

Using SSL for Secure Connections

To establish an SSL connection, set the SSL information in the connectionData dictionary. For example:

```
mysql-js> var mysqlx=require('mysqlx').mysqlx;

mysql-js> var session=mysqlx.getNodeSession({host: 'localhost',
                                             dbUser: 'root',
                                             dbPassword: 'mypasswd',
                                             ssl_ca: "path_to_ca_file",
                                             ssl_cert: "path_to_cert_file",
                                             ssl_key: "path_to_key_file"});

mysql-js> print(session)
<NodeSession:root@localhost>
mysql-js>
```

19.2.2.7 MySQL Shell Global Variables

MySQL Shell reserves certain variables as global variables, which are assigned to commonly used objects in scripting. This section describes the available global variables and provides examples of working with them. The global variables are:

- `session` represents the global session if one has been established.
- `db` represents a schema if one has been defined, for example by a URI.

MySQL Shell provides interactive error resolution for common situations related to using the global variables. For example:

- Attempting to use an undefined `session` global variable.
- Attempting to retrieve an nonexistent schema using `session`.
- Attempting to use an undefined `db` global variable.

Undefined Global Session

The global `session` variable is set when a global session is established. When a global session is established, issuing a `session` statement in MySQL Shell displays the session type and its URI as follows:

```
mysql-js> session
<XSession:root@localhost:33060>
mysql-js>
```

If no global session has been established, MySQL Shell displays the following:

```
mysql-js> session
<Undefined>
mysql-js>
```

If you attempt to use the `session` variable when no global session is established, interactive error resolution starts and you are prompted to provide the required information to establish a global session. If the session is successfully established, it is assigned to the `session` variable. The prompts are:

- An initial prompt explains that no global session is established and asks if one should be established.
- If the you choose to set a global session, the session type is requested.
- Either the URI or the alias of a stored session is requested.
- If required, a password is requested.

For example:

```
mysql-js> session.uri
The global session is not set, do you want to establish a session? [y/N]: y
Please specify the session type:

1) X
2) Node
3) Classic

Type: 2
Please specify the MySQL server URI (or $slias): root@localhost
Enter password:*******
root@localhost:33060
mysql-js>
```

Undefined db Variable

The global `db` variable is set when a global session is established and a default schema is configured. For example, using a URI such as `root@localhost/sakila` to establish a global session connected to the MySQL Server at `localhost`, on port 33060, as the user `root`, assigns the schema `sakila` to the global variable `db`. Once a schema is defined, issuing `db` at the MySQL Shell prompt prints the schema name as follows:

```
mysql-js> db
<Schema:world_x>
mysql-js>
```

If there is no global session established, the following is displayed:

```
mysql-js> db
<Undefined>
mysql-js>
```

If you attempt to use the `db` variable when no global session has been established, the following error is displayed:

```
mysql-js> db.getCollections()
LogicError: The db variable is not set, establish a global session first.
at (shell):1:2
in db.getCollections()
^
```

If a global session has been established but you attempt to use an undefined `db`, interactive error resolution begins and you are prompted to define an active schema by providing the schema name. If this succeeds the `db` variable is set to the defined schema. For example:

```
mysql-js> db.getCollections()
The db variable is not set, do you want to set the active schema? [y/N]:y
Please specify the schema:world_x
[
<Collection:countryinfo>
]
mysql-js> db
<Schema:world_x>
```

```
mysql-js>
```

Retrieving an Nonexistent Schema

If you attempt to use `session` to retrieve an nonexistent schema, interactive error resolution provides the option to create the schema.

```
mysql-js> var mySchema = session.getSchema('my_test')
The schema my_test does not exist, do you want to create it? [y/N]: y

mysql-js> mySchema
<Schema:my_test>
mysql-js>
```

In all cases, if you do not provide the information required to resolve each situation, a proper result of executing the requested statement on an undefined variable is displayed.

19.3 MySQL Shell Code Execution

This section explains how code execution works in MySQL Shell.

19.3.1 Interactive Code Execution

The default mode of MySQL Shell provides interactive execution of database operations that you type at the command prompt. These operations can be written in JavaScript, Python or SQL depending on the type of session being used. When executed, the results of the operation are displayed on-screen.

As with any other language interpreter, MySQL Shell is very strict regarding syntax. For example, the following JavaScript snippet reads and prints the documents in a collection:

```
var mysqlx = require('mysqlx').mysqlx;
var mySession = mysqlx.getSession('user:pwd@localhost');
var result = mySession.world_x.countryinfo.find().execute();
var record = result.fetchOne();
while(record){
  print(record);
  record = result.fetchOne();
}
```

As seen above, the call to `find()` is followed by the `execute()` function. CRUD database commands are only actually executed on the MySQL Server when `execute()` is called. However, when working with MySQL Shell interactively, `execute()` is implicitly called whenever you press `Return` on a statement. Then the results of the operation are fetched and displayed on-screen. The rules for when you need to call `execute()` or not are as follows:

- When using MySQL Shell in this way, calling `execute()` becomes optional on:
  - `Collection.add()`
  - `Collection.find()`
  - `Collection.remove()`
  - `Collection.modify()`
  - `Table.insert()`
  - `Table.select()`
  - `Table.delete()`
  - `Table.update()`

    - `NodeSession.sql()`

- Automatic execution is disabled if the object is assigned to a variable. In such a case calling `execute()` is mandatory to perform the operation.
- When a line is processed and the function returns any of the available `Result` objects, the information contained in the Result object is automatically displayed on screen. The functions that return a Result object include:
    - The SQL execution and CRUD operations (listed above)
    - Transaction handling and drop functions of the session objects in both mysql and mysqlx modules: -
        - `startTransaction()`
        - `commit()`
        - `rollback()`
        - `dropSchema()`
        - `dropTable()`
        - `dropCollection()`
        - `dropView()`
        - `ClassicSession.runSql()`

Based on the above rules, the statements needed in the MySQL Shell in interactive mode to establish a session, query, and print the documents in a collection are:

```
mysql-js> var mysqlx = require('mysqlx').mysqlx;
mysql-js> var mySession = mysqlx.getSession('user:pwd@localhost');
```

No call to `execute()` is needed and the Result object is automatically printed.

```
mysql-js> mySession.world_x.countryinfo.find();
```

19.3.2 Batch Code Execution

As well as interactive code execution, MySQL Shell provides batch code execution from:

- A file loaded for processing.
- A file containing code that is redirected to the standard input for execution.
- Code from a different source that is redirected to the standard input for execution.

For example:

Loading SQL code from a file for batch processing.

```
shell> mysqlsh --file code.js
```

Redirecting a file to standard input for execution.

```
shell> mysqlsh < code.js
```

Redirecting code to standard input for execution.

```
shell> echo "show databases;" | mysqlsh --sql --uri root@192.168.1.141:33060
```

Executable Scripts

Starting with version 1.0.4, on Linux you can create executable scripts that run with MySQL Shell by including a `#!` line as the first line of the script. This line should provide the full path to MySQL Shell and include the `--file` option. For example:

```
#!/usr/local/mysql-shell/bin/mysqlsh --file
print("Hello World\n");
```

The script file must be marked as executable in the filesystem. Running the script invokes MySQL Shell and it executes the contents of the script.

19.3.3 Output Formats

The output of the commands processed on the server can be formatted in different ways. This section details the different available output formats.

19.3.3.1 Table Format

The table format is used by default when MySQL Shell is in interactive mode. The output is presented as a formatted table for a better view and to aid analysis.

```
mysql-sql> select * from sakila.actor limit 3;
+----------+-------------+----------------+---------------------+
| actor_id | first_name  | last_name      | last_update         |
+----------+-------------+----------------+---------------------+
|        1 | PENELOPE    | GUINESS        | 2006-02-15 4:34:33  |
|        2 | NICK        | WAHLBERG       | 2006-02-15 4:34:33  |
|        3 | ED          | CHASE          | 2006-02-15 4:34:33  |
+----------+-------------+----------------+---------------------+
3 rows in set (0.00 sec)

mysql-sql>
```

To get this output format when running in batch mode, use the `--table` command-line option.

19.3.3.2 Tab Separated Format

This format is used by default when running MySQL Shell in batch mode, to have better output for automated analysis.

```
>echo "select * from sakila.actor limit 3;" | mysqlsh --classic --uri root@192.168.1.141:33460
actor_id        first_name      last_name       last_update
1       PENELOPE        GUINESS 2006-02-15 4:34:33
2       NICK    WAHLBERG        2006-02-15 4:34:33
3       ED      CHASE   2006-02-15 4:34:33
```

19.3.3.3 JSON Format Output

MySQL Shell supports the JSON format for output and it is available both in interactive and batch mode. This output format can be enabled using the `--json` command-line option:

JSON Format in Batch Mode

```
shell>echo "select * from sakila.actor limit 3;" | mysqlsh --json --sqlc --uri root@192.168.1.141:3306
{"duration":"0.00 sec","info":"","row_count":3,"rows":[[1,"PENELOPE","GUINESS",{"year":2006,"month":1,"day"
```

```
shell>echo "select * from sakila.actor limit 3;" | mysqlsh --json=raw --sqlc --uri root@192.168.1.141:3
{"duration":"0.00 sec","info":"","row_count":3,"rows":[[1,"PENELOPE","GUINESS",{"year":2006,"month":1,"

shell>echo "select * from sakila.actor limit 3;" | mysqlsh --json=pretty --sqlc --uri root@192.168.1.14
{
    "duration": "0.00 sec",
    "info": "",
    "row_count": 3,
    "rows": [
        [
            1,
            "PENELOPE",
            "GUINESS",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ],
        [
            2,
            "NICK",
            "WAHLBERG",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ],
        [
            3,
            "ED",
            "CHASE",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ]
    ],
    "warning_count": 0
}
shell>
```

JSON Format in Interactive Mode (started with --json=raw)

```
mysql-sql> select * from sakila.actor limit 3;
{"duration":"0.00 sec","info":"","row_count":3,"rows":[[1,"PENELOPE","GUINESS",{"year":2006,"month":1,"

mysql-sql>
```

JSON Format in Interactive Mode (started with --json=pretty)

```
mysql-sql> select * from sakila.actor limit 3;
{
    "duration": "0.00 sec",
    "info": "",
    "row_count": 3,
    "rows": [
        [
```

```
            1,
            "PENELOPE",
            "GUINESS",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ],
        [
            2,
            "NICK",
            "WAHLBERG",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ],
        [
            3,
            "ED",
            "CHASE",
            {
                "year": 2006,
                "month": 1,
                "day": 15,
                "hour": 4,
                "minute": 34,
                "second": 33.0
            }
        ]
    ],
    "warning_count": 0
}

mysql-sql>
```

19.3.3.4 Result Metadata

When an operation is executed, in addition to any results returned, some additional information is available. This includes information such as the number of affected rows, warnings, duration, and so on, when any of these conditions is true:

- JSON format is being used for the output
- MySQL Shell is running in interactive mode.

19.3.4 Active Language

MySQL Shell can execute SQL, JavaScript or Python code, but only one language can be active at a time. The active mode determines how the executed statements are processed:

- If using SQL mode, statements are processed as SQL which means they are sent to the MySQL server for execution.
- If using JavaScript mode, statements are processed as JavaScript code.
- If using Python mode, statements are processed as Python code.

When running MySQL Shell in interactive mode, activate a specific language by entering the commands: `\sql`, `\js`, `\py`.

When running MySQL Shell in batch mode, activate a specific language by passing any of these command-line options: `--js`, `--py` or `--sql`. The default mode if none is specified is JavaScript.

Use MySQL Shell to execute the content of the file `code.sql` as SQL.

```
shell> mysqlsh --sql < code.sql
```

Use MySQL Shell to execute the content of the file `code.js` as JavaScript code.

```
shell> mysqlsh < code.js
```

Use MySQL Shell to execute the content of the file `code.py` as Python code.

```
shell> mysqlsh --py < code.py
```

19.3.5 Batch Mode Made Interactive

This section describes code execution in batch mode.

- In batch mode, all the command logic described above is not available, only valid code for the active language can be executed.
- When processing SQL code, it is executed statement by statement using the following logic: read/process/print result.
- When processing non-SQL code, it is loaded entirely from the input source and executed as a unit.

Use the `--interactive` (or `-i`) command-line option to configure MySQL Shell to process the input source as if it were being issued in interactive mode; this enables all the features provided by the Interactive mode to be used in batch processing.

Note

In this case, whatever the source is, it is read line by line and processed using the interactive pipeline.

19.3.5.1 Multiple-line Support

It is possible to specify statements over multiple lines. When in Python or JavaScript mode, multiple-line mode is automatically enabled when a block of statements starts like in function definitions, if/then statements, for loops, and so on. In SQL mode multiple line mode starts when the command `\` is issued.

Once multiple-line mode is started, the subsequently entered statements are cached.

For example:

```
mysql-sql> \
... create procedure get_actors()
... begin
...   select first_name from sakila.actor;
... end
...
mysql-sql>
```

19.4 Configuring MySQL Shell

This section explains how to configure MySQL Shell using commands executable from the interactive code editor and command-line options. For a description of MySQL Shell command-line options, see `mysqlsh` — The MySQL Shell.

19.4.1 MySQL Shell Commands

MySQL Shell provides commands which enable you to modify the execution environment of the code editor, for example to configure the active programming language or a MySQL Server connection. The following table lists the commands that are available regardless of the currently selected language. As commands need to be available independent of the *execution mode*, they start with an escape sequence, the `\` character.

Command	Alias/Shortcut	Description
`\help`	`\h` or `\?`	Prints help about MySQL Shell commands.
`\quit`	`\q` or `\exit`	Exit MySQL Shell.
`\`		In SQL mode, begin multiple-line mode. Code is cached and executed when an empty line is entered.
`\status`		Show the current MySQL Shell status.
`\js`		Switch execution mode to JavaScript.
`\py`		Switch execution mode to Python.
`\sql`		Switch execution mode to SQL.
`\connect`	`\c`	Connect to a MySQL Server with a URI using an XSession (X Protocol).
`\connect_node`	`\cn`	(Removed in version 1.0.4, use `\connect -n`) Connect to a MySQL Server with a URI using a Node session.
`\connect_classic`	`\cc`	(Removed in version 1.0.4, use `\connect -c`) Connect to a MySQL Server with a URI using a Classic session (MySQL Protocol).
`\use`		Specify the schema to use.
`\source`	`\.`	Execute a script file using the active language.
`\warnings`	`\W`	Show any warnings generated by a statement.
`\nowarnings`	`\w`	Do not show any warnings generated by a statement.
`\lsconn`	`\lsc`	Print the connection data for the stored sessions.
`\saveconn`	`\savec`	Save connection data of a session, optionaly use `-f` to force overwriting an existing connection.
`\addconn`	`\addc`	(Removed in version 1.0.4, see `\saveconn`) Store the connection data of a session.
`\rmconn`		Removes a stored session.
`\chconn`		(Removed in version 1.0.4, see `\saveconn`) Updates a stored session.

Help Command

The `\help` command can be used with or without parameters. When used without parameters a general help is printed including information about:

- Available commands.
- Available commands for the active mode.

When used with a parameter, the parameter must be a valid command. If that is the case, help for that specific command is printed including:

- Description
- Supported aliases if any

- Additional help if any

For example:

```
\help connect
```

If the parameter is not a valid command, the general help is printed.

Connect Command

The `\connect` command is used to connect to a MySQL Server using an URI. This command creates an X Protocol connection by default.

For example:

```
\connect root@localhost:3306
```

If a password is required you are prompted for it.

Use the `-n` option to create a Node session, using the X Protocol to connect to a single server. For example:

```
\connect -n root@localhost:3306
```

Use the `-c` option to create a Classic session, enabling you to use the MySQL Protocol to issue SQL commands directly on a server. For example:

```
\connect -c root@localhost:3306
```

Source Command

The `\source` command is used to execute code from a script at a given path. For example:

```
\source /tmp/mydata.sql
```

It can be used to execute either SQL, JavaScript or Python code. The code in the file is executed using the active language, so to process SQL code the MySQL Shell must be in SQL mode.

Warning

As the code is executed using the active language, executing a script in a different language than the currently selected execution mode language could lead to unexpected results.

Use Command

The `\use` command enables you to choose which schema is active, for example:

```
\use schema_name
```

The `\use` command requires a global development session to be active. If the global development session is an XSession then the `\use` command only sets `db` to the object representing the *`schema_name`* but does not set a current schema on the database. If the global development session is one of NodeSession or ClassicSession the `\use` command sets the current schema to the specified *`schema_name`* and updates the `db` variable to the object that represents the selected schema.

19.5 Stored Sessions

This section describes how MySQL Shell stores a persistent list of connection data for sessions, identified by a session name.

There are two approaches to manage and use stored connection data:

- Through MySQL Shell commands.
- Through a built-in object available for JavaScript and Python.

With these options you can add, update, remove, list and use stored connection data.

The connections are stored as a text file in the MySQL Shell folder, usually found in the user directory.

The file name is `stored_sessions.json`

In Windows a typical location is:

```
%APPDATA%\MySQL\mysqlsh\stored_sessions.json
```

In Unix a typical location is:

```
~/.mysqlsh/stored_sessions.json
```

19.5.1 MySQL Shell Stored Session Commands

The following commands are used to work with stored sessions within MySQL Shell.

Adding a Connection

To add connection data for a session use:

```
\saveconn name [URI]
```

- *`name`* : Mandatory parameter, the name that is assigned to the stored connection data. Note that this must be a valid identifier:
  - Starts with an alphabetic character.
  - Only contains alphanumeric characters and underscore.
- *`URI`* : Optional parameter containing connection data in URI format. If not specified and a global session is active, the URI for the global session is used instead.
- `-f` : Optional parameter to force overwriting an existing connection. If an attempt is made to store a connection, but the provided *`name`* already exists, an error occurs. Use this parameter to force storing the connection, which overrides what was previously stored under the given session name.

Updating a Stored Connection

To update the connection data for a specific session name use:

```
\saveconn -f name URI
```

- *`name`* : Mandatory parameter, indicates the connection to be updated.
- `URI` : Mandatory parameter, contains the connection data that replaces the connection for the indicated session name.

Removing a Stored Connection

To remove a connection use:

```
\rmconn name
```

- *name* : Mandatory parameter that indicates the connection data to be removed.

Listing the Stored Connections

To list the stored connections use:

```
\lsconn
```

Using a Stored Connection

Stored connections can be used to create a global session by issuing a `\connect name` command. The stored session is loaded and used to create the global session.

When using a stored connection through the connect shell commands, you are prompted for the password if it is missing from the stored connection data.

19.5.2 Scripting Stored Sessions

A built in object is available to manage and use the stored connections through JavaScript or Python: `shell.storedSessions`

19.5.2.1 Adding a Connection

To add a new connection use:

```
shell.storedSessions.add(name, connectionData[, override])
```

- *name* : Mandatory parameter, it is the session name that identifies the connection data. This must be a valid identifier.
- `connectionData` : Mandatory parameter, it contains the connection information and can be either a string in URI format or a connection data map containing any of the following values:
  - host
  - port
  - dbUser
  - dbPassword
  - schema
  - ssl_ca
  - ssl_cert
  - ssl_key

  To store the connection data at least the host and dbUser parameter should be specified, the rest are optional.
- **override** : Optional boolean parameter, used to indicate the desired behavior in case there is a stored connection data associated to the same name. Default value is false.
  - If override is false and connection data with the same name already exists, an error is generated.
  - In the same scenario, if the new connection data overrides the previous data and an error is not generated.

19.5.2.2 Updating a Stored Connection

To update the connection data for a specific session name use:

```
shell.storedSessions.update(name, connectionData)
```

- **name** : Mandatory parameter, it is the name that will identify the connection data. Note that this must be a valid identifier.
- **connectionData** : Mandatory parameter, it contains the connection data and can be either a string in URI format or a connection data map.

19.5.2.3 Removing a Stored Connection

To remove a connection use:

```
shell.storedSessions.remove(name)
```

- **name** : Mandatory parameter, identifies the connection data to be removed.

19.5.2.4 Listing the Stored Connections

To list the stored connections use:

```
shell.storedSessions
```

19.5.2.5 Using a Stored Connection

When the `shell.storedSessions` object is initialized, it loads the stored session data and creates a property to access them as `shell.storedSessions.`*`name`*

```
var session = mysqlx.getSession(shell.storedSessions.name)
```

When used this way, if the connection data has no password it is not requested. It is the developers responsibility to provide the password either using the stored session data or as an additional parameter as follows:

```
var session = mysqlx.getSession(shell.storedSessions.name, password)
```

19.6 MySQL Shell Application Log

This section explains the application log.

19.6.1 Application Log

MySQL Shell can be configured to generate an application log file with information about issues of varying severity. You can use this information to verify the state of MySQL Shell while it is running.

The location of the log file is the user configuration path and the file is named `mysqlsh.log`.

Log File on Windows

On Windows, the default path to the log file is `%APPDATA%\MySQL\mysqlsh\mysqlsh.log`

To find the location of `%APPDATA%` on your system, echo it from the comand-line. For example:

```
```

```
C:>echo %APPDATA%

C:\Users\exampleuser\AppData\Roaming
```

On Windows, the path is determined by the result of gathering the `%APPDATA%` folder specific to that user, and then appending `MySQL\mysqlsh`. Using the above example, we end up with:

```
C:\Users\exampleuser\AppData\Roaming\MySQL\mysqlsh\mysqlsh.log
```

Log File on Unix-based Systems

For a machine running Unix, the default path is `~/.mysqlsh/mysqlsh.log` where "~" represents the user's home directory. The environment variable `HOME` also represents the user's home directory. Appending `.mysqlsh` to the user's home directory determines the default path to the logs. For example:

```
C:>echo $HOME

/home/exampleuser

shell> less /home/exampleuser/.mysqlsh/mysqlsh.log
```

These paths can be overridden on all platforms by defining the environment variable `MYSQL_USER_CONFIG_PATH`. The value of this variable replaces `%APPDATA%` in Windows or `$HOME` in Unix.

By default, logging is disabled in MySQL Shell. To enable logging use the `--log-level` command-line option when starting MySQL Shell. For example:

```
shell> mysqlsh --log-level=4
```

The number assigned to `--log-level` controls the level of detail in the log. The following logging levels are supported:

Log Level	Meaning
1	None, the default
2	Internal Error
3	Error
4	Warning
5	Info
6	Debug
7	Debug2
8	Debug3

The log format is plain text and entries contain a timestamp and description of the problem, along with the log level from the above list. For example:

```
2016-04-05 22:23:01: Error: Default Domain: (shell):1:8: MySQLError: You have an error
in your SQL syntax; check the manual that corresponds to your MySQL server version for
the right syntax to use near '' at line 1 (1064) in session.sql("select * from t
limit").execute().all();
```

19.7 Customizing MySQL Shell

MySQL Shell offers the ability to customize the behavior and code execution environment through startup scripts, which are executed when the application is first run. Using such scripts enables you to:

- Add additional search paths for Python or JavaScript modules.
- Override the default prompt used by the Python and JavaScript modes.
- Define global functions or variables.
- Any other possible initialization through JavaScript or Python.

19.7.1 Working With Start-Up Scripts

When MySQL Shell enters either into JavaScript or Python mode, it searches for startup scripts to be executed. The startup scripts are JavaScript or Python specific scripts containing the instructions to be executed when the corresponding mode is initialized.

Startup scripts must be named as follows:

- For JavaScript mode: `mysqlshrc.js`
- For Python mode: `mysqlshrc.py`

MySQL Shell searches the following paths for these files (in order of execution).

On Windows:

1. `%PROGRAMDATA%MySQLmysqlshmysqlshrc.[js|py]`
2. `%MYSQLSH_HOME%sharedmysqlshmysqlshrc.[js|py]`
3. `<mysqlsh binary path>mysqlshrc.[js|py]`
4. `%APPDATA%MySQLmysqlshmysqlshrc.[js|py]`

On Linux and OSX:

1. `/etc/mysql/mysqlsh/mysqlshrc.[js|py]`
2. `$MYSQLSH_HOME/shared/mysqlsh/mysqlshrc.[js|py]`
3. `<mysqlsh binary path>/mysqlshrc.[js|py]`
4. `$HOME/.mysqlsh/mysqlshrc.[js|py]`

The environment variable `MYSQLSH_HOME` defines the root folder of a standard setup of MySQL Shell. If `MYSQLSH_HOME` is not defined it is automatically calculated based on the location of the MySQL Shell binary, therefore on many standard setups it is not required to define `MYSQLSH_HOME`.

If `MYSQLSH_HOME` is not defined and the MySQL Shell binary is not in a standard install folder structure, then the path defined in option 3 in the above lists is used. If using a standard install or if `MYSQLSH_HOME` points to a standard install folder structure, then the path defined in option 3 is not used.

Warning

The lists above also define the order of searching the paths, so if something is defined in two different scripts, the script executed later takes precedence.

19.7.2 Adding Module Search Paths

There are two ways to add additional module search paths:

- Through environment variables
- Through startup scripts

19.7.2.1 Environment Variables

Python uses the `PYTHONPATH` environment variable to allow extending the search paths for python modules. The value of this variable is a list of paths separated by:

- A colon character in Linux and OSX
- A semicolon character in Windows

To achieve this in JavaScript, MySQL Shell supports defining additional JavaScript module paths using the `MYSQLSH_JS_MODULE_PATH` environment variable. The value of this variable is a list of semicolon separated paths.

19.7.2.2 Startup Scripts

The addition of module search paths can be achieved for both languages through the corresponding startup script.

For Python modify the `mysqlshrc.py` file and append the required paths into the `sys.path` array.

```
# Import the sys module
import sys

# Append the additional module paths
sys.path.append('~/custom/python')
sys.path.append('~/other/custom/modules')
```

For JavaScript the same task is achieved by adding code into the `mysqlshrc.js` file to append the required paths into the predefined `shell.js_module_paths` array.

```
// Append the additional module paths
shell.js.module_paths[shell.js.module_paths.length] = '~/custom/js';
shell.js.module_paths[shell.js.module_paths.length] = '~/other/custom/modules';
```

19.7.3 Overriding the Default Prompt

MySQL Shell uses a default prompt for both Python (`mysql-py>`) and JavaScript (`mysql-js>`).

You can customize the language specific prompt using the `shell.custom_prompt()` function. This function must return a string that is used as the prompt. To have a custom prompt when MySQL Shell starts, define this function in a startup script. The following example shows how this functionality can be used.

In Python `shell.custom_prompt()` could be defined as:

```
# Import the sys module
from time import gmtime, strftime

def my_prompt():
  ret_val = strftime("%H:%M:%S", gmtime())

  if session and session.isOpen():
    data = shell.parseUri(session.getUri())

    ret_val = "%s-%s-%s-py> " % (ret_val, data.dbUser, data.host)
  else:
    ret_val = "%s-disconnected-py> " % ret_val

  return ret_val


shell.custom_prompt = my_prompt
```

In JavaScript `shell.custom_prompt()` could be defined as:

```
shell.custom_prompt = function(){
  var now = new Date();

  var ret_val = now.getHours().toString()+ ":" + now.getMinutes().toString() + ":" + now.getSeconds().toStr

  if (session && session.isOpen()){
    var data = shell.parseUri(session.getUri());

    ret_val += "-" + data.dbUser + "-" + data.host + "-js> ";
  }
  else
    ret_val += "-disconnected-js> ";

  return ret_val;
}
```

The following example demonstrates using the custom prompt functions defined above in startup script. The prompts show the current system time, and if a session is open the current user and host:

```
Welcome to MySQL Shell 1.0.4 Development Preview

Copyright (c) 2016, Oracle and/or its affiliates. All rights reserved.

Oracle is a registered trademark of Oracle Corporation and/or its
affiliates. Other names may be trademarks of their respective
owners.

Type '\help', '\h' or '\?' for help.

Currently in JavaScript mode. Use \sql to switch to SQL mode and execute queries.
14:34:32-disconnected-js> \py
Switching to Python mode...

19:34:39-disconnected-py> \connect root:@localhost
Creating an X Session to root@localhost:33060
No default schema selected.

19:34:50-root-localhost-py> \js
Switching to JavaScript mode...
14:34:57-root-localhost-js>
```

Chapter 20 Using MySQL as a Document Store

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This chapter introduces an alternative way of working with MySQL as a document store, sometimes referred to as “using NoSQL”. If your intention is to use MySQL in a traditional (SQL) way, this chapter is probably not relevant to you.

Using MySQL as a document store is currently a preproduction feature to which this notice applies: Section 20.1, “Preproduction Status — Legal Notice”.

Relational databases such as MySQL usually required a document schema to be defined before documents can be stored. The features described in this section enable you to use MySQL as a document store, which is a schema-less, and therefore schema-flexible, storage system for documents. When using MySQL as a document store, to create documents describing products you do not need to know and define all possible attributes of any products before storing them and operating with them. This differs from working with a relational database and storing products in a table, when all columns of the table must be known and defined before adding any products to the database. The features described in this chapter enable you to choose how you configure MySQL, using only the document store model, or combining the flexibility of the document store model with the power of the relational model.

These sections cover the usage of MySQL as a document store:

- The Section 20.2, “Key Concepts” section covers concepts like Document, Collection, Session, and Schema to help you understand how to use MySQL as a document store.
- The Section 20.3, “Setting Up MySQL as a Document Store” section explains how to configure X Plugin on a MySQL Server, so it can function as a document store, and how to install MySQL Shell to use as a client.

- The MySQL Shell is an interactive interface to MySQL supporting JavaScript, Python, or SQL modes. You can use the MySQL Shell to prototype applications, execute queries and update data. The *quick-start guides (tutorials)* help you to get started using MySQL Shell.

  The quick-start guide for JavaScript is here: Section 20.4, “Quick-Start Guide: MySQL Shell for JavaScript”.

  The quick-start guide for Python is here: Section 20.5, “Quick-Start Guide: MySQL Shell for Python”.

- Chapter 19, *MySQL Shell User Guide* provides more detailed information about using MySQL Shell.

- *X DevAPI User guide.*

  Clients that communicate with a MySQL Server using the X Protocol can use the X DevAPI to develop applications. For example MySQL Shell and MySQL Connectors provide this ability by implementing the X DevAPI. X DevAPI offers a modern programming interface with a simple yet powerful design which provides support for established industry standard concepts. See X DevAPI User Guide for in-depth tutorials on using X DevAPI.

- The following MySQL Connectors support the X Protocol and enable you to use X DevAPI in your chosen language to develop applications that communicate with a MySQL Server functioning as a document store.

  - MySQL Shell X DevAPI and AdminAPI reference: JavaScript and Python.
  - MySQL Connector/J X DevAPI Reference
  - MySQL Connector/Node.js X DevAPI Reference
  - MySQL Connector/Net X DevAPI Reference
  - MySQL Connector/Python X DevAPI Reference

  Each implementation of X DevAPI provides library reference documentation, available at the above links.

20.1 Preproduction Status — Legal Notice

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20.2 Key Concepts

This section explains the concepts introduced as part of using MySQL as a document store.

Document

A Document is a set of key and value pairs, as represented by a JSON object. A Document is represented internally using the MySQL binary JSON object, through the JSON MySQL datatype. The values of fields can contain other documents, arrays, and lists of documents.

```
{
    "GNP": .6,
    "IndepYear": 1967,
    "Name": "Sealand",
    "_id": "SEA",
    "demographics": {
        "LifeExpectancy": 79,
```

```
        "Population": 27
    },
    "geography": {
        "Continent": "Europe",
        "Region": "British Islands",
        "SurfaceArea": 193
    },
    "government": {
        "GovernmentForm": "Monarchy",
        "HeadOfState": "Michael Bates"
    }
}
```

Collection

A Collection is a container that may be used to store Documents in a MySQL database.

CRUD Operations

Create, Read, Update and Delete (CRUD) operations are the four basic operations that can be performed on a database Collection or Table. In terms of MySQL this means:

- Create a new entry (insertion or addition)
- Read entries (queries)
- Update entries
- Delete entries

X Plugin

The MySQL Server plugin which enables communication using X Protocol. Supports clients that implement X DevAPI and enables you to use MySQL as a document store.

X Protocol

A protocol to communicate with a MySQL Server running X Plugin. X Protocol supports both CRUD and SQL operations, authentication via SASL, allows streaming (pipelining) of commands and is extensible on the protocol and the message layer.

20.3 Setting Up MySQL as a Document Store

This section describes how to install the X Plugin to enable MySQL Server to use the X Protocol and use MySQL Shell to act as a client to the server.

The prerequisite for using MySQL as a document store is the X Plugin, which implements the X Protocol. Without the X Plugin running, X Protocol clients cannot connect to the server. The X Plugin is shipped with MySQL — installing it does not involve a separate download.

Setting up MySQL as a document store follows the steps outlined here:

1. **Install MySQL 8 or upgrade from MySQL 5.7.12 or lower.**

   When the installation or upgrade is done, start the server. For server startup instructions, see Section 2.9.2, “Starting the Server”.

Note

MySQL Installer enables you to perform this and the next step (Install the X Plugin) at the same time for new installations on Microsoft Windows. In the **Plugin and Extensions** screen, check mark the **Enable X Protocol/**

> **MySQL as a Document Store** check box. After the installation, verify that the X Plugin has been installed.

2. **Install the X Plugin.** A non-root account can be used to install the plugin as long as the account has `INSERT` privilege for the `mysql.plugin` table.

   Always save your existing configuration settings before reconfiguring the server.

   To install the built-in X Plugin, do one of the following:

   - Using MySQL Installer for Windows:

     a. Launch MySQL Installer for Windows. MySQL Installer dashboard opens.

     b. Click the **Reconfigure** quick action for MySQL Server. Use **Next** and **Back** to configure the following items:

        - In **Accounts and Roles**, confirm the current `root` account password.
        - In **Plugin and Extensions**, check mark the **Enable X Protocol/MySQL as a Document Store** check box. MySQL Installer provides a default port number and opens the firewall port for network access.
        - In **Apply Server Configuration**, click **Execute**.
        - Click **Finish** to close MySQL Installer.

     c. Install MySQL Shell.

   - Using MySQL Shell:

     a. Install MySQL Shell.

     b. Open a terminal window (command prompt on Windows) and navigate to the MySQL binaries location (for example, `/usr/bin/` on Linux).

     c. Run the following command:

     ```
     mysqlsh -u user -h localhost --classic --dba enableXProtocol
     ```

   - Using the MySQL Client program:

     a. Open a terminal window (command prompt on Windows) and navigate to the MySQL binaries location (for example, `/usr/bin/` on Linux).

     b. Invoke the `mysql` command-line client:

     ```
     mysql -u user -p
     ```

     c. Issue the following statement:

     ```
     mysql> INSTALL PLUGIN mysqlx SONAME 'mysqlx.so';
     ```

     Replace `mysqlx.so` with `mysqlx.dll` for Windows.

     d. Install MySQL Shell.

3. **Verify that the X Plugin has been installed.**

   When the X Plugin is installed properly, it shows up in the list when you query for active plugins on the server with one of the following commands:

- MySQL Shell command:

```
mysqlsh -u user --sqlc -e "show plugins"
```

- MySQL Client program command:

```
mysql -u user -p -e "show plugins"
```

If you encounter problems with the X Plugin installation, or if you want to learn about alternative ways of installing, configuring, or uninstalling server plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”.

mysqlxsys@localhost User Account

Installing the X Plugin creates a mysqlxsys@localhost user account. If, for some reason, creating the user account fails, the X Plugin installation fails, too. Here is an explanation on what the mysqlxsys@localhost user account is for and what to do when its creation fails.

The X Plugin installation process uses the MySQL root user to create an internal account for the mysqlxsys@localhost user. The mysqlxsys@localhost account is used by the X Plugin for authentication of external users against the MySQL account system and for killing sessions when requested by a privileged user. The mysqlxsys@localhost account is created as locked, so it cannot be used to log in by external users. If for some reason the MySQL root account is not available, before you start the X Plugin installation you must manually create the mysqlxsys@localhost user by issuing the following statements in the mysql command-line client:

```
CREATE USER IF NOT EXISTS mysqlxsys@localhost IDENTIFIED WITH
mysql_native_password AS 'password' ACCOUNT LOCK;
GRANT SELECT ON mysql.user TO mysqlxsys@localhost;
GRANT SUPER ON *.* TO mysqlxsys@localhost;
```

Uninstalling the X Plugin

If you ever want to uninstall (deactivate) the X Plugin, issue the following statement in the mysql command-line client:

```
UNINSTALL PLUGIN mysqlx;
```

Do not use MySQL Shell to issue the previous statement. It works from MySQL Shell, but you get an error (code 1130). Also, uninstalling the plugin removes the mysqlxsys user.

20.3.1 Installing MySQL Shell

Important

For the Community version of MySQL Shell: Before installing MySQL Shell, make sure you have the Visual C++ Redistributable for Visual Studio 2015 (available at the Microsoft Download Center) installed on your Windows system.

Note

MySQL Shell 1.0.8 is not supplied with an MSI Installer. See Installing MySQL Shell Binaries for the manual install procedure.

This section describes how to download, install, and start MySQL Shell, which is an interactive JavaScript, Python, or SQL interface supporting development and administration for the MySQL Server. MySQL Shell is a component that you can install separately.

Requirements

MySQL Shell is available on Microsoft Windows, Linux, and OS X for 64-bit platforms. MySQL Shell requires that the built-in X Plugin be active. You can install the server plugin before or after you install MySQL Shell. For instructions, see Installing the X Plugin.

20.3.1.1 Installing MySQL Shell on Microsoft Windows

Important

The Community version of MySQL Shell requires the Visual C++ Redistributable for Visual Studio 2013 to work; make sure that is installed on your Windows system before installing MySQL Shell.

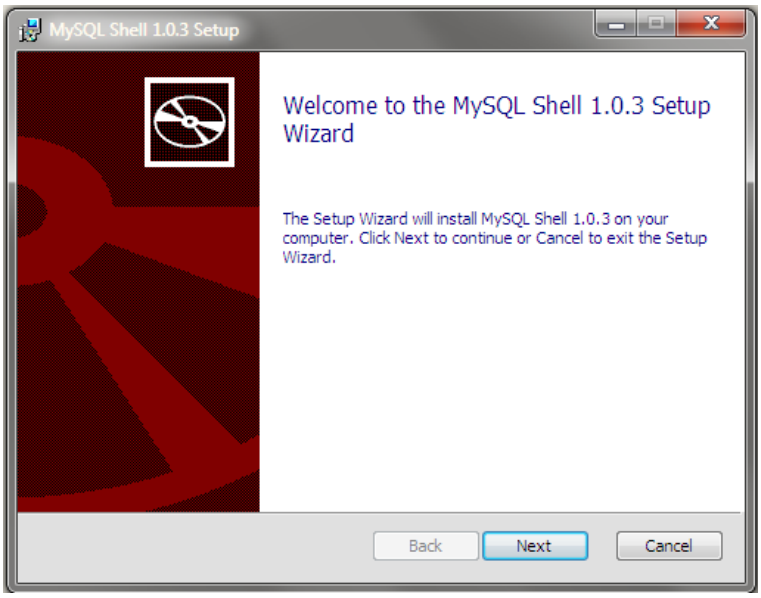
Note

MySQL Shell is currently not supplied with an MSI Installer. See Installing MySQL Shell Binaries for the manual install procedure.

To install MySQL Shell on Microsoft Windows using the MSI Installer, do the following:

1. Download the **Windows (x86, 64-bit), MSI Installer** package from http://dev.mysql.com/downloads/shell/.
2. When prompted, click **Run**.
3. Follow the steps in the Setup Wizard.

Figure 20.1 Installation of MySQL Shell on Windows



If you have installed MySQL without enabling the X Plugin, then later on decide you want to install the X Plugin, or if you are installing MySQL *without* using MySQL Installer, see Installing the X Plugin.

20.3.1.2 Installing MySQL Shell on Linux

Note

Installation packages for MySQL Shell are available only for a limited number of Linux distributions, and only for 64-bit systems.

For supported Linux distributions, the easiest way to install MySQL Shell on Linux is to use the MySQL APT repository or MySQL Yum repository. For systems not using the MySQL repositories, MySQL Shell can also be downloaded and installed directly.

Installing MySQL Shell with the MySQL APT Repository

Note

Installation of MySQL Shell using the MySQL APT repository is only supported on Ubuntu 14.04 LTS (“Trusty Tahr”) and Ubuntu 15.10 (“Wily Werewolf”).

For Linux distributions supported by the MySQL APT repository, follow one of the paths below:

- If you do not yet have the MySQL APT repository as a software repository on your system, do the following:
  - Follow the steps given in Adding the MySQL APT Repository, paying special attention to the following:
    - During the installation of the configuration package, when asked in the dialogue box to configure the repository, make sure you choose MySQL 5.7 (which is the default option) as the release series you want, and enable the **MySQL Preview Packages** component.
    - Make sure you do not skip the step for updating package information for the MySQL APT repository:

      ```
      sudo apt-get update
      ```
  - Install MySQL Shell with this command:

    ```
    sudo apt-get install mysql-shell
    ```
- If you already have the MySQL APT repository as a software repository on your system, do the following:
  - Update package information for the MySQL APT repository:

    ```
    sudo apt-get update
    ```
  - Update the MySQL APT repository configuration package with the following command:

    ```
    sudo apt-get install mysql-apt-config
    ```

    When asked in the dialogue box to configure the repository, make sure you choose MySQL 5.7 (which is the default option) as the release series you want, and enable the **MySQL Preview Packages** component.
  - Install MySQL Shell with this command:

    ```
    sudo apt-get install mysql-shell
    ```

Installing MySQL Shell with the MySQL Yum Repository

Note

Installation of MySQL Shell using the MySQL Yum repository is only supported on EL6 and EL7-based systems, as well as Fedora 23 and 24.

For Linux distributions supported by the MySQL Yum repository, follow these steps to install MySQL Shell:

- Do one of the following:
  - If you already have the MySQL Yum repository as a software repository on your system and the repository was configured with the new release package `mysql57-community-release`, skip to the next step (“Enable the MySQL Tools Preview subrepository...”).
  - If you already have the MySQL Yum repository as a software repository on your system but have configured the repository with the old release package `mysql-community-release`, it is easiest to install MySQL Shell by first reconfiguring the MySQL Yum repository with the new `mysql57-community-release` package. To do so, you need to remove your old release package first, with the following command :

    ```
    sudo yum remove mysql-community-release
    ```

    For dnf-enabled systems, do this instead:

    ```
    sudo dnf erase mysql-community-release
    ```

    Then, follow the steps given in Adding the MySQL Yum Repository to install the new release package, `mysql57-community-release`.
  - If you do not yet have the MySQL Yum repository as a software repository on your system, follow the steps given in Adding the MySQL Yum Repository.
- Enable the MySQL Tools Preview subrepository. You can do that by editing manually the `/etc/yum.repos.d/mysql-community.repo` file. This is an example of the subrepository's default entry in the file (the `baseurl` entry in your file might look different, depending on your Linux distribution):

  ```
  [mysql-tools-preview]
  name=MySQL Tools Preview
  baseurl=http://repo.mysql.com/yum/mysql-tools-preview/el/6/$basearch/
  enabled=0
  gpgcheck=1
  gpgkey=file:///etc/pki/rpm-gpg/RPM-GPG-KEY-mysql
  ```

  Change the entry `enabled=0` to `enabled=1` to enable the subrepository.
- Install MySQL Shell with this command:

  ```
  sudo yum install mysql-shell
  ```

  For dnf-enabled systems, do this instead:

  ```
  sudo dnf install mysql-shell
  ```

Installing MySQL Shell from Direct Downloads from the MySQL Developer Zone

RPM, Debian, and source packages for installing MySQL Shell are also available for download at Download MySQL Shell.

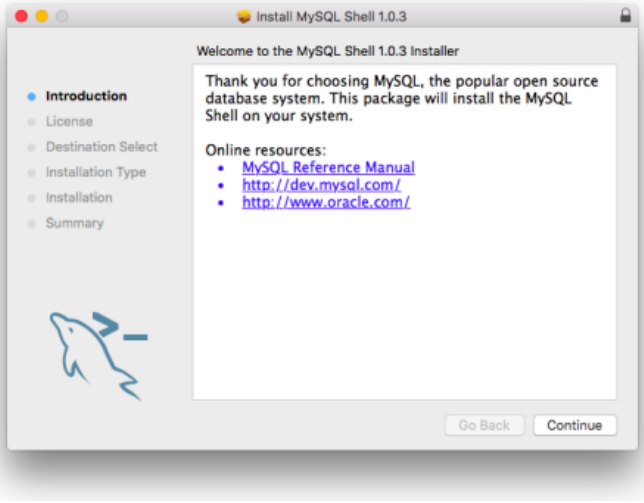
20.3.1.3 Installing MySQL Shell on OS X

To install MySQL Shell on OS X, do the following:

1. Download the package from http://dev.mysql.com/downloads/shell/.
2. Double-click the downloaded DMG to mount it. Finder opens.
3. Double-click the `.pkg` file shown in the Finder window.

4. Follow the steps in the installation wizard.

Figure 20.2 Installation of MySQL Shell on OS X



5. When the installer finishes, eject the DMG. (It can be deleted.)

20.3.2 Starting MySQL Shell

You need an account name and password to establish a session using MySQL Shell. Replace *`user`* with your account name.

On the same system where the server instance is running, open a terminal window (command prompt on Windows) and start MySQL Shell with the following command:

```
mysqlsh --uri user@localhost
```

You are prompted to input your password and then this establishes an X Session.

For instructions to get you started using MySQL Shell and MySQL as a document store, see the following quick-start guides:

- Quick-Start Guide: MySQL Shell for JavaScript
- Quick-Start Guide: MySQL Shell for Python

20.4 Quick-Start Guide: MySQL Shell for JavaScript

This quick-start guide provides instructions to begin prototyping database applications interactively with MySQL Shell. The guide includes the following topics:

- Introduction to MySQL functionality, MySQL Shell, and the `world_x` database sample.
- Operations to manage collections and documents.
- Operations to manage relational tables.
- Operations that apply to documents within tables.

Available Quick-Start Guides

- MySQL Shell for Python

- MySQL for Visual Studio

Related Information

- Chapter 19, *MySQL Shell User Guide* provides more in-depth information about MySQL Shell.
- X DevAPI User Guide provides more examples of using X DevAPI.

20.4.1 Introduction

The MySQL Shell for JavaScript quick start provides a short but comprehensive introduction to the database functionality including the new X DevAPI, which offers a modern, integrative way to work with relational and document data, without requiring SQL knowledge from application developers.

In MySQL, tables are the native data storage container type and collections are stored internally using tables.

JSON Documents and Collections

A JSON document is a data structure composed of field/value pairs stored within a collection. The values of fields often contain other documents, arrays, and lists of documents.

```
{
    "GNP": .6,
    "IndepYear": 1967,
    "Name": "Sealand",
    "_id": "SEA",
    "demographics": {
        "LifeExpectancy": 79,
        "Population": 27
    },
    "geography": {
        "Continent": "Europe",
        "Region": "British Islands",
        "SurfaceArea": 193
    },
    "government": {
        "GovernmentForm": "Monarchy",
        "HeadOfState": "Michael Bates"
    }
}
```

Relational Tables

A table in MySQL enables you to store data organized in rows and columns. The structure of a table is defined by one or more columns with user-defined names and data types. Every row stored in the table has the same structure.

```
+------+------------------+-------------+-----------------+------------------------+
| ID   | Name             | CountryCode | District        | Info                   |
+------+------------------+-------------+-----------------+------------------------+
|    1 | Kabul            | AFG         | Kabol           |{"Population": 1780000} |
|    2 | Qandahar         | AFG         | Qandahar        |{"Population": 237500}  |
|    3 | Herat            | AFG         | Herat           |{"Population": 186800}  |
|    4 | Mazar-e-Sharif   | AFG         | Balkh           |{"Population": 127800}  |
|    5 | Amsterdam        | NLD         | Noord-Holland   |{"Population": 731200}  |
|    6 | Rotterdam        | NLD         | Zuid-Holland    |{"Population": 593321}  |
+------+------------------+-------------+-----------------+------------------------+
```

Related Information

- Chapter 19, *MySQL Shell User Guide* provides a general overview.

- See X DevAPI User Guide for development reference documentation.

20.4.2 Import Database Sample

The `world_x` database sample contains one JSON collection and a set of three relational tables:

- Collection
  - countryinfo: Information about countries in the world.
- Tables
  - country: Minimal information about countries of the world.
  - city: Information about some of the cities in those countries.
  - countrylanguage: Languages spoken in each country.

Requirements

You must install MySQL Shell with the X Protocol enabled. For instructions, see Section 20.3, “Setting Up MySQL as a Document Store”.

Start the server before you load the `world_x` database for this guide.

Download world_x Database

To prepare the `world_x` database sample, follow these steps:

1. Download world_x-db.zip.
2. Extract the installation archive to a temporary location such as `/tmp/`. Unpacking the archive results in a single file named `world_x.sql`.
3. Create the schema with the following command:

```
mysqlsh -u root --sql < /tmp/world_x-db/world_x.sql
Enter password: ****
Records: 4079  Duplicates: 0  Warnings: 0

Records: 239  Duplicates: 0  Warnings: 0

Records: 239  Duplicates: 0  Warnings: 0

Records: 984  Duplicates: 0  Warnings: 0
```

   Enter your password when prompted. A non-root account can be used as long as the account has privileges to create new databases.

   Replace `/tmp/` with the path to the `world_x.sql` file on your system.

Related Information

- Section 19.2.1, “MySQL Shell Sessions” explains session types.
- See Chapter 2, *Installing and Upgrading MySQL* for general installation assistance.

20.4.3 MySQL Shell

MySQL Shell is a unified scripting interface to MySQL Server. It supports scripting in JavaScript and Python. JavaScript is the default processing mode. In most cases, you need an account to connect to the local MySQL server instance.

Start MySQL Shell

After you have installed and started MySQL server, connect MySQL Shell to the server instance. By default, MySQL Shell connects using the X Protocol.

On the same system where the server instance is running, open a terminal window and start MySQL Shell with the following command:

```
mysqlsh name@localhost/world_x
Creating a Session to 'root@localhost/world_x'
Enter password: ****
```

You may need to specify the path as appropriate.

In addition:

- *name* represents the user name of your MySQL account.
- MySQL Shell prompts you for your password.
- The default schema for this session is the `world_x` database. For instructions on setting up the `world_x` database sample, see Section 20.4.2, “Import Database Sample”.

The `mysql-js>` prompt indicates that the active language for this session is JavaScript.

```
mysql-js>
```

When you run `mysqlsh` without the host argument, MySQL Shell attempts to connect to the server instance running on the localhost interface on port 33060. To specify a different host or port number, as well as other options, see the option descriptions at `mysqlsh` — The MySQL Shell.

MySQL Shell supports input-line editing as follows:

- **left-arrow** and **right-arrow** keys move horizontally within the current input line.
- **up-arrow** and **down-arrow** keys move up and down through the set of previously entered lines.
- **Backspace** deletes the character before the cursor and typing new characters enters them at the cursor position.
- **Enter** enters the current input line.

Get Help for MySQL Shell

Type `mysqlsh --help` at the prompt of your command interpreter for a list of command-line options.

```
mysqlsh --help
```

Type `\help` at the MySQL Shell prompt for a list of available commands and their descriptions.

```
mysql-js> \help
```

Type `\help` followed by a command name for detailed help about an individual MySQL Shell command. For example, to view help on the `\connect` command, type:

```
mysql-js> \help \connect
```

Quit MySQL Shell

To quit MySQL Shell, type the following command:

```
mysql-js> \quit
```

Related Information

- See Section 19.3.1, “Interactive Code Execution” for an explanation of how interactive code execution works in MySQL Shell.
- See Section 19.2, “Getting Started with MySQL Shell” to learn about session and connection alternatives.

20.4.4 Documents and Collections

In MySQL, collections contain JSON documents that you can add, find, update, and remove. Collections are containers within a schema that you create, list, and drop.

The examples in this section use the countryinfo collection in the `world_x` database. For instructions on setting up the `world_x` database sample, see Section 20.4.2, “Import Database Sample”.

Documents

In MySQL, documents are represented as JSON objects. Internally, they are stored in an efficient binary format that enables fast lookups and updates.

- Simple document format for JavaScript:

```
{field1: "value", field2 : 10, "field 3": null}
```

An array of documents consists of a set of documents separated by commas and enclosed within `[` and `]` characters.

- Simple array of documents for JavaScript:

```
[{Name: "Aruba", _id: "ABW"}, {Name: "Angola", _id: "AGO"}]
```

MySQL supports the following JavaScript value types in JSON documents:

- numbers (integer and floating point)
- strings
- boolean (false and true)
- null
- arrays of more JSON values
- nested (or embedded) objects of more JSON values

Collections

Collections are containers for documents that share a purpose and possibly share one or more indexes. Each collection has a unique name and exists within a single schema.

The term schema is equivalent to a database, which means a group of database objects (as opposed to relational schema used to enforce structure and constraints over data). A schema does not enforce conformity on the documents in a collection.

In this quick-start guide:

- Basic objects include:

Object form	Description
`db`	`db` is a global variable assigned to the current active schema that you specified on the command line. You can type `db` in MySQL Shell to print a description of the object, which in this case will be the name of the schema it represents.
`db.getCollections()`	db.getCollections() holds a list of collections in the schema. Use the list to get references to collection objects, iterate over them, and so on.

- Basic operations scoped by collections include:

Operation form	Description
`db.name.add()`	The add() method inserts one document or a list of documents into the named collection.
`db.name.find()`	The find() method returns some or all documents in the named collection.
`db.name.modify()`	The modify() method updates documents in the named collection.
`db.name.remove()`	The remove() method deletes one document or a list of documents from the named collection.

Related Information

- See Working with Collections for a general overview.
- CRUD EBNF Definitions provides a complete list of operations.

20.4.4.1 Create, List, and Drop Collections

In MySQL Shell, you can create new collections, get a list of the existing collections in a schema, and remove an existing collection from a schema. Collection names are case-sensitive and each collection name must be unique.

Confirm the Schema

To show the value that is assigned to the schema variable, type `db`.

```
mysql-js> db
<Schema:world_x>
```

If the schema value is not `Schema:world_x`, then set the `db` variable as follows:

```
mysql-js> \use world_x
Schema `world_x` accessible through db.
```

Create a Collection

To create a new collection in an existing schema, use the `createCollection()` method.

The following example creates in the `world_x` database a collection called flags. The method returns a collection object.

```
mysql-js> db.createCollection("flags")
<Collection:flags>
```

List Collections

To display all collections in the `world_x` database, use the `getCollections()` method on the schema object. Collections returned by the server appear between brackets.

```
mysql-js> db.getCollections()
[
    <Collection:countryinfo>,
    <Collection:flags>
]
```

Drop a Collection

To drop an existing collection from a database, use the `dropCollection()` method on the session object. For example, to drop the flags collection from the `world_x` database, type:

```
mysql-js> session.dropCollection("world_x", "flags")
Query OK (0.04 sec)
```

Related Information

- See Section 19.2.2.6, “Connections in JavaScript and Python” to learn more about the session object.
- See Collection Objects for more examples.

20.4.4.2 Add Documents

You can use the `add()` method to insert one document or a list documents into an existing collection using MySQL Shell. All examples in this section use the countryinfo collection.

Confirm the Schema

To show the value that is assigned to the schema variable, type `db`.

```
mysql-js> db
<Schema:world_x>
```

If the schema value is not `Schema:world_x`, then set the `db` variable as follows:

```
mysql-js> \use world_x
Schema `world_x` accessible through db.
```

Add a Document

Insert the following document into the countryinfo collection. Press **Enter** twice to insert the document.

```
mysql-js> db.countryinfo.add(
 {
    GNP: .6,
    IndepYear: 1967,
    Name: "Sealand",
    _id: "SEA",
    demographics: {
        LifeExpectancy: 79,
        Population: 27
    },
    geography: {
        Continent: "Europe",
        Region: "British Islands",
        SurfaceArea: 193
```

```
    },
    government: {
        GovernmentForm: "Monarchy",
        HeadOfState: "Michael Bates"
    }
  }
)
Query OK, 1 item affected (0.02 sec)
```

The method returns the status of the operation.

Each document requires an identifier field called `_id`. The value of the `_id` field must be unique among all documents in the same collection. If the document passed to the `add()` method does not contain the `_id` field, MySQL Shell automatically inserts a field into the document and sets the value to a generated universal unique identifier (UUID).

Related Information

- See CollectionAddFunction for the full syntax definition.

20.4.4.3 Find Documents

You can use the `find()` method to query for and return documents from a collection in a database. MySQL Shell provides additional methods to use with the `find()` method to filter and sort the returned documents.

MySQL provides the following operators to specify search conditions: `OR` (`||`), `AND` (`&&`), `XOR`, `IS`, `NOT`, `BETWEEN`, `IN`, `LIKE`, `!=`, `<>`, `>`, `>=`, `<`, `<=`, `&`, `|`, `<<`, `>>`, `+`, `-`, `*`, `/`, `~`, and `%`.

Find All Documents in a Collection

To return all documents in a collection, use the `find()` method without specifying search conditions. For example, the following operation returns all documents in the countryinfo collection.

```
mysql-js> db.countryinfo.find()
[
     {
          "GNP": 828,
          "IndepYear": null,
          "Name": "Aruba",
          "_id": "ABW",
          "demographics": {
              "LifeExpectancy": 78.4000015258789,
              "Population": 103000
          },
          "geography": {
              "Continent": "North America",
              "Region": "Caribbean",
              "SurfaceArea": 193
          },
          "government": {
              "GovernmentForm": "Nonmetropolitan Territory of The Netherlands",
              "HeadOfState": "Beatrix"
          }
          ...
      }
 ]
240 documents in set (0.00 sec)
```

The method produces results that contain operational information in addition to all documents in the collection.

An empty set (no matching documents) returns the following information:

```
Empty set (0.00 sec)
```

Filter Searches

You can include search conditions with the `find()` method. The syntax for expressions that form a search condition is the same as that of traditional MySQL. You must enclose all expressions in quotes.

All examples in this section use the countryinfo collection in the `world_x` database. For the sake of brevity, some of the examples do not display output.

A simple search condition consists of the `_id` field and unique identifier of a document. The following example returns a single document matching the identifier string:

```
mysql-js> db.countryinfo.find("_id = 'AUS'")
[
    {
        "GNP": 351182,
        "IndepYear": 1901,
        "Name": "Australia",
        "_id": "AUS",
        "demographics": {
            "LifeExpectancy": 79.80000305175781,
            "Population": 18886000
        },
        "geography": {
            "Continent": "Oceania",
            "Region": "Australia and New Zealand",
            "SurfaceArea": 7741220
        },
        "government": {
            "GovernmentForm": "Constitutional Monarchy, Federation",
            "HeadOfState": "Elisabeth II"
        }
    }
]
1 document in set (0.01 sec)
```

The following example searches for all countries that have a GNP higher than $500 billion. The countryinfo collection measures GNP in units of million.

```
mysql-js> db.countryinfo.find("GNP > 500000")
...[output removed]
10 documents in set (0.00 sec)
```

The Population field in the following query is embedded within the demographics object. To access the embedded field, use a period between demographics and Population to identify the relationship. Document and field names are case sensitive.

```
mysql-js> db.countryinfo.find("GNP > 500000 and demographics.Population < 100000000")
...[output removed]
6 documents in set (0.00 sec)
```

Arithmetic operators in the following expression are used to query for countries with a GNP per capita higher than $30000. Search conditions can include arithmetic operators and most MySQL functions.

Note

Seven documents in the countryinfo collection have a population value of zero. Warning messages appear at the end of the output.

```
mysql-js> db.countryinfo.find("GNP*1000000/demographics.Population > 30000")
...[output removed]
9 documents in set, 7 warnings (0.00 sec)
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
```

```
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
```

You can separate a value from the search condition by using the `bind()` method. For example, instead of specifying a hard-coded country name as the condition, substitute a named placeholder consisting of a colon followed by a name that begins with a letter, such as *country*. Then include the placeholder and value in the `bind()` method as follows:

```
mysql-js> db.countryinfo.find("Name = :country").bind("country", "Italy")
[
    {
        "GNP": 1161755,
        "IndepYear": 1861,
        "Name": "Italy",
        "_id": "ITA",
        "demographics": {
            "LifeExpectancy": 79,
            "Population": 57680000
        },
        "geography": {
            "Continent": "Europe",
            "Region": "Southern Europe",
            "SurfaceArea": 301316
        },
        "government": {
            "GovernmentForm": "Republic",
            "HeadOfState": "Carlo Azeglio Ciampi"
        }
    }
]
1 document in set (0.01 sec)
```

Tip

Within a program, binding enables you to specify placeholders in your expressions, which are filled in with values before execution and can benefit from automatic escaping, as appropriate.

Always use binding to sanitize input. Avoid introducing values in queries using string concatenation, which can produce invalid input and, in some cases, can cause security issues.

Project Results

You can return specific fields of a document, instead of returning all the fields. The following example returns the GNP and Name fields of all documents in the countryinfo collection matching the search conditions.

Use the `fields()` method to pass the list of fields to return.

```
mysql-js> db.countryinfo.find("GNP > 5000000").fields(["GNP", "Name"])
[
    {
        "GNP": 8510700,
        "Name": "United States"
    }
]
1 document in set (0.00 sec)
```

In addition, you can alter the returned documents—adding, renaming, nesting and even computing new field values—with an expression that describes the document to return. For example, alter the names of the fields with the following expression to return only two documents.

```
mysql-js> db.countryinfo.find().
fields(mysqlx.expr('{"Name": upper(Name), "GNPPerCapita": GNP*1000000/demographics.Population}')).
limit(2)
[
    {
        "GNPPerCapita": 8038.834951456311,
        "Name": "ARUBA"
    },
    {
        "GNPPerCapita": 263.0281690140845,
        "Name": "AFGHANISTAN"
    }
]
2 documents in set (0.00 sec)
```

Limit, Sort, and Skip Results

You can apply the `limit()`, `sort()`, and `skip()` methods to manage the number and order of documents returned by the `find()` method.

To specify the number of documents included in a result set, append the `limit()` method with a value to the `find()` method. The following query returns the first five documents in the countryinfo collection.

```
mysql-js> db.countryinfo.find().limit(5)
... [output removed]
5 documents in set (0.00 sec)
```

To specify an order for the results, append the `sort()` method to the `find()` method. Pass to the `sort()` method a list of one or more fields to sort by and, optionally, the descending (`desc`) or ascending (`asc`) attribute as appropriate. Ascending order is the default order type.

For example, the following query sorts all documents by the IndepYear field and then returns the first eight documents in descending order.

```
mysql-js> db.countryinfo.find().sort(["IndepYear desc"]).limit(8)
... [output removed]
8 documents in set (0.00 sec)
```

By default, the `limit()` method starts from the first document in the collection. You can use the `skip()` method to change the starting document. For example, to ignore the first document and return the next eight documents matching the condition, pass to the `skip()` method a value of 1.

```
mysql-js> db.countryinfo.find().sort(["IndepYear desc"]).limit(8).skip(1)
... [output removed]
8 documents in set (0.00 sec)
```

Related Information

- The MySQL Reference Manual provides detailed documentation on functions and operators.
- See CollectionFindFunction for the full syntax definition.

20.4.4.4 Modify Documents

You can use the `modify()` method to update one or more documents in a collection. The X DevAPI provides additional methods for use with the `modify()` method to:

- Set and unset fields within documents.
- Append, insert, and delete arrays.

- Bind, limit, and sort the documents to be modified.

Set and Unset Fields

The `modify()` method works by filtering a collection to include only the documents to be modified and then applying the operations that you specify to those documents.

In the following example, the `modify()` method uses the search condition to identify the document to change and then the `set()` method replaces two values within the nested demographics object.

```
mysql-js> db.countryinfo.modify("_id = 'SEA'").
set("demographics", {LifeExpectancy: 78, Population: 28})
Query OK, 1 item affected (0.04 sec)
```

After you modify a document, use the `find()` method to verify the change.

To remove content from a document, use the `modify()` and `unset()` methods. For example, the following query removes the GNP from a document that matches the search condition.

```
mysql-js> db.countryinfo.modify("Name = 'Sealand'").unset("GNP")
Query OK, 1 item affected (0.01 sec)
```

Use the `find()` method to verify the change.

```
mysql-js> db.countryinfo.find("Name = 'Sealand'")
[
    {
        "IndepYear": 1967,
        "Name": "Sealand",
        "_id": "SEA",
        "demographics": {
            "LifeExpectancy": 78,
            "Population": 28
        },
        "geography": {
            "Continent": "Europe",
            "Region": "British Islands",
            "SurfaceArea": 193
        },
        "government": {
            "GovernmentForm": "Monarchy",
            "HeadOfState": "Michael Bates"
        }
    }
]
1 document in set (0.00 sec)
```

Append, Insert, and Delete Arrays

To append an element to an array field, or insert, or delete elements in an array, use the `arrayAppend()`, `arrayInsert()`, or `arrayDelete()` methods. The following examples modify the countryinfo collection to enable tracking of international airports.

The first example uses the `modify()` and `set()` methods to create a new Airports field in all documents.

Caution

Use care when you modify documents without specifying a search condition. This action will modify all documents in the collection.

```
mysql-js> db.countryinfo.modify().set("Airports", [])
Query OK, 240 items affected (0.07 sec)
```

With the Airports field added, the next example uses the `arrayAppend()` method to add a new airport to one of the documents. *$.Airports* in the following example represents the Airports field of the current document.

```
mysql-js> db.countryinfo.modify("Name = 'France'").arrayAppend("$.Airports", "ORY")
Query OK, 1 item affected (0.02 sec)
```

Use **`db.countryinfo.find("Name = 'France'")`** to see the change.

To insert an element at a different position in the array, use the `arrayInsert()` method to specify which index to insert in the path expression. In this case, the index is 0, or the first element in the array.

```
mysql-js> db.countryinfo.modify("Name = 'France'").arrayInsert("$.Airports[0]", "CDG")
Query OK, 1 item affected (0.04 sec)
```

To delete an element from the array, you must pass to the `arrayDelete()` method the index of the element to be deleted.

```
mysql-js> db.countryinfo.modify("Name = 'France'").arrayDelete("$.Airports[1]")
Query OK, 1 item affected (0.03 sec)
```

Related Information

- The MySQL Reference Manual provides instructions to help you search for and modify JSON values.
- See CollectionModifyFunction for the full syntax definition.

20.4.4.5 Remove Documents

You can use the `remove()` method to delete some or all documents from a collection in a database. The X DevAPI provides additional methods for use with the `remove()` method to filter and sort the documents to be removed.

Remove Documents Using Conditions

The example that follows passes a search condition to the `remove()` method. All documents matching the condition will be removed from the countryinfo collection. In this example, one document matches the condition.

```
mysql-js> db.countryinfo.remove("_id = 'SEA'")
Query OK, 1 item affected (0.02 sec)
```

Remove the First Document

To remove the first document in the countryinfo collection, use the `limit()` method with a value of 1.

```
mysql-js> db.countryinfo.remove().limit(1)
Query OK, 1 item affected (0.03 sec)
```

Remove the Last Document in an Order

The following example removes the last document in the countryinfo collection by country name.

```
mysql-js> db.countryinfo.remove().sort(["Name desc"]).limit(1)
Query OK, 1 item affected (0.02 sec)
```

Remove All Documents in a Collection

You can remove all documents in a collection. To do so, use the `remove()` method without specifying a search condition.

Caution

Use care when you remove documents without specifying a search condition. This action will delete all documents in the collection.

Related Information

- See CollectionRemoveFunction for the full syntax definition.
- See Section 20.4.2, “Import Database Sample” for instructions to recreate the `world_x` database.

20.4.4.6 Create and Drop Indexes

Indexes are used to find documents with specific field values quickly. Without an index, MySQL must begin with the first document and then read through the entire collection to find the relevant fields. The larger the collection, the more this costs. If a collection is large and queries on a specific field are common, then consider creating an index on a specific field inside a document.

For example, the following query will perform better with an index:

```
mysql-js> db.countryinfo.find("demographics.Population < 100")
...[output removed]
8 documents in set (0.00 sec)
```

The `createIndex()` method creates an index that you can define as nonunique or unique. Use the `field()` method to chain the fields that should be indexed. The `execute()` method is required to create or drop an index.

In MySQL, the `_id` field is equivalent to a primary key by default.

Add a Nonunique Index

To create a nonunique index, pass to the `createIndex()` method an index name. Duplicate index names are prohibited.

In the following example, the first parameter of the `field()` method specifies the Population field inside the demographics object and the next parameter indicates that the field should be indexed as an Integer numeric value. The last parameter indicates whether the field should require the NOT NULL constraint. If the value is `false`, the field can contain `NULL` values.

```
mysql-js> db.countryinfo.createIndex("pop").
field("demographics.Population", "INTEGER", false).execute()
Query OK (0.04 sec)
```

Add a Unique Index

To create a unique index, pass to the `createIndex()` method an index name and the `mysqlx.IndexType.UNIQUE` type. Country `"Name"` is another common field in the countryinfo collection to index. In the following example, `"Text(40)"` represents the number of characters to index and `true` indicates that the field cannot contain any `NULL` values.

```
mysql-js> db.countryinfo.createIndex("name", mysqlx.IndexType.UNIQUE).
field("Name", "TEXT(40)", true).execute()
Query OK (0.04 sec)
```

Drop an Index

To drop an index, pass to the `dropIndex()` method the name of the index to drop. For example, you can drop the “pop” index as follows:

```
mysql-js> db.countryinfo.dropIndex("pop").execute()
Query OK (0.58 sec)
```

Related Information

- See Collection Index Management Functions for the full syntax definition.

20.4.5 Relational Tables

You can use MySQL Shell to manipulate not just JSON documents, but also relational tables.

In MySQL, each relational table is associated with a particular storage engine. The examples in this section use `InnoDB` tables in the `world_x` database.

Confirm the Schema

To show the value that is assigned to the schema variable, type `db`.

```
mysql-js> db
<Schema:world_x>
```

If the schema value is not `Schema:world_x`, then set the `db` variable as follows:

```
mysql-js> \use world_x
Schema `world_x` accessible through db.
```

Show All Tables

To display all relational tables in the `world_x` database, use the `getTables()` method on the schema object.

```
mysql-js> db.getTables()
{
    "city": <Table:city>,
    "country": <Table:country>,
    "countrylanguage": <Table:countrylanguage>
}
```

Basic Table Operations

Basic operations scoped by tables include:

Operation form	Description
`db.name.insert()`	The insert() method inserts one or more records into the named table.
`db.name.select()`	The select() method returns some or all records in the named table.
`db.name.update()`	The update() method updates records in the named table.
`db.name.delete()`	The delete() method deletes one or more records from the named table.

Related Information

- See Working with Relational Tables for a general overview.
- CRUD EBNF Definitions provides a complete list of operations.
- See Section 20.4.2, “Import Database Sample” for instructions on setting up the `world_x` database sample.

20.4.5.1 Insert Records into Tables

You can use the `insert()` method with the `values()` method to insert records into an existing relational table. The `insert()` method accepts individual columns or all columns in the table. Use one or more `values()` methods to specify the values to be inserted.

Insert a Complete Record

To insert a complete record, pass to the `insert()` method all columns in the table. Then pass to the `values()` method one value for each column in the table. For example, to add a new record to the city table in the `world_x` database, insert the following record and press **Enter** twice.

```
mysql-js> db.city.insert("ID", "Name", "CountryCode", "District", "Info").
values(null, "Olympia", "USA", "Washington", '{"Population": 5000}')
Query OK, 1 item affected (0.01 sec)
```

The city table has five columns: ID, Name, CountryCode, District, and Info. Each value must match the data type of the column it represents.

Insert a Partial Record

The following example inserts values into the ID, Name, and CountryCode columns of the city table.

```
mysql-js> db.city.insert("ID", "Name", "CountryCode").
values(null, "Little Falls", "USA").values(null, "Happy Valley", "USA")
Query OK, 2 item affected (0.03 sec)
```

When you specify columns using the `insert()` method, the number of values must match the number of columns. In the previous example, you must supply three values to match the three columns specified.

Related Information

- See TableInsertFunction for the full syntax definition.

20.4.5.2 Select Tables

You can use the `select()` method to query for and return records from a table in a database. The X DevAPI provides additional methods to use with the `select()` method to filter and sort the returned records.

MySQL provides the following operators to specify search conditions: `OR` (`||`), `AND` (`&&`), `XOR`, `IS`, `NOT`, `BETWEEN`, `IN`, `LIKE`, `!=`, `<>`, `>`, `>=`, `<`, `<=`, `&`, `|`, `<<`, `>>`, `+`, `-`, `*`, `/`, `~`, and `%`.

Select All Records

To issue a query that returns all records from an existing table, use the `select()` method without specifying search conditions. The following example selects all records from the city table in the `world_x` database.

Note

Limit the use of the empty `select()` method to interactive statements. Always use explicit column-name selections in your application code.

```
mysql-js> db.city.select()
+------+------------+-------------+------------+-------------------------+
| ID   | Name       | CountryCode | District   | Info                    |
+------+------------+-------------+------------+-------------------------+
|    1 | Kabul      | AFG         | Kabol      |{"Population": 1780000}  |
|    2 | Qandahar   | AFG         | Qandahar   |{"Population": 237500}   |
|    3 | Herat      | AFG         | Herat      |{"Population": 186800}   |
```

```
...    ...          ...           ...          ...
| 4079 | Rafah      | PSE         | Rafah      |{"Population": 92020}    |
+------+------- ----+-------------+------------+-------------------------+
4082 rows in set (0.01 sec)
```

An empty set (no matching records) returns the following information:

```
Empty set (0.00 sec)
```

Filter Searches

To issue a query that returns a set of table columns, use the `select()` method and specify the columns to return between square brackets. This query returns the Name and CountryCode columns from the city table.

```
mysql-js> db.city.select(["Name", "CountryCode"])
+-------------------+-------------+
| Name              | CountryCode |
+-------------------+-------------+
| Kabul             | AFG         |
| Qandahar          | AFG         |
| Herat             | AFG         |
| Mazar-e-Sharif    | AFG         |
| Amsterdam         | NLD         |
...                 ...
| Rafah             | PSE         |
| Olympia           | USA         |
| Little Falls      | USA         |
| Happy Valley      | USA         |
+-------------------+-------------+
4082 rows in set (0.00 sec)
```

To issue a query that returns rows matching specific search conditions, use the `where()` method to include those conditions. For example, the following example returns the names and country codes of the cities that start with the letter Z.

```
mysql-js> db.city.select(["Name", "CountryCode"]).where("Name like 'Z%'")
+-------------------+-------------+
| Name              | CountryCode |
+-------------------+-------------+
| Zaanstad          | NLD         |
| Zoetermeer        | NLD         |
| Zwolle            | NLD         |
| Zenica            | BIH         |
| Zagazig           | EGY         |
| Zaragoza          | ESP         |
| Zamboanga         | PHL         |
| Zahedan           | IRN         |
| Zanjan            | IRN         |
| Zabol             | IRN         |
| Zama              | JPN         |
| Zhezqazghan       | KAZ         |
| Zhengzhou         | CHN         |
...                 ...
| Zeleznogorsk      | RUS         |
+-------------------+-------------+
59 rows in set (0.00 sec)
```

You can separate a value from the search condition by using the `bind()` method. For example, instead of using "Name = 'Z%' " as the condition, substitute a named placeholder consisting of a colon followed by a name that begins with a letter, such as *name*. Then include the placeholder and value in the `bind()` method as follows:

```
mysql-js> db.city.select(["Name", "CountryCode"]).
              where("Name like :name").bind("name", "Z%")
```

Tip

Within a program, binding enables you to specify placeholders in your expressions, which are filled in with values before execution and can benefit from automatic escaping, as appropriate.

Always use binding to sanitize input. Avoid introducing values in queries using string concatenation, which can produce invalid input and, in some cases, can cause security issues.

Project Results

To issue a query using the `AND` operator, add the operator between search conditions in the `where()` method.

```
mysql-js> db.city.select(["Name", "CountryCode"]).
              where("Name like 'Z%' and CountryCode = 'CHN'")
+---------------+-------------+
| Name          | CountryCode |
+---------------+-------------+
| Zhengzhou     | CHN         |
| Zibo          | CHN         |
| Zhangjiakou   | CHN         |
| Zhuzhou       | CHN         |
| Zhangjiang    | CHN         |
| Zigong        | CHN         |
| Zaozhuang     | CHN         |
...             ...
| Zhangjiagang  | CHN         |
+---------------+-------------+
22 rows in set (0.01 sec)
```

To specify multiple conditional operators, you can enclose the search conditions in parenthesis to change the operator precedence. The following example demonstrates the placement of `AND` and `OR` operators.

```
mysql-js> db.city.select(["Name", "CountryCode"]).
where("Name like 'Z%' and (CountryCode = 'CHN' or CountryCode = 'RUS')")
+------------------+-------------+
| Name             | CountryCode |
+------------------+-------------+
| Zhengzhou        | CHN         |
| Zibo             | CHN         |
| Zhangjiakou      | CHN         |
| Zhuzhou          | CHN         |
...                ...
| Zeleznogorsk     | RUS         |
+------------------+-------------+
29 rows in set (0.01 sec)
```

Limit, Order, and Offset Results

You can apply the `limit()`, `orderBy()`, and `offSet()` methods to manage the number and order of records returned by the `select()` method.

To specify the number of records included in a result set, append the `limit()` method with a value to the `select()` method. For example, the following query returns the first five records in the country table.

```
mysql-js> db.country.select(["Code", "Name"]).limit(5)
+------+-------------+
| Code | Name        |
+------+-------------+
| ABW  | Aruba       |
| AFG  | Afghanistan |
```

```
| AGO  | Angola      |
| AIA  | Anguilla    |
| ALB  | Albania     |
+------+-------------+
5 rows in set (0.00 sec)
```

To specify an order for the results, append the `orderBy()` method to the `select()` method. Pass to the `orderBy()` method a list of one or more columns to sort by and, optionally, the descending (`desc`) or ascending (`asc`) attribute as appropriate. Ascending order is the default order type.

For example, the following query sorts all records by the Name column and then returns the first three records in descending order .

```
mysql-js> db.country.select(["Code", "Name"]).orderBy(["Name desc"]).limit(3)
+------+------------+
| Code | Name       |
+------+------------+
| ZWE  | Zimbabwe   |
| ZMB  | Zambia     |
| YUG  | Yugoslavia |
+------+------------+
3 rows in set (0.00 sec)
```

By default, the `limit()` method starts from the first record in the table. You can use the `offset()` method to change the starting record. For example, to ignore the first record and return the next three records matching the condition, pass to the `offset()` method a value of 1.

```
mysql-js> db.country.select(["Code", "Name"]).orderBy(["Name desc"]).limit(3).offset(1)
+------+------------+
| Code | Name       |
+------+------------+
| ZMB  | Zambia     |
| YUG  | Yugoslavia |
| YEM  | Yemen      |
+------+------------+
3 rows in set (0.00 sec)
```

Related Information

- The MySQL Reference Manual provides detailed documentation on functions and operators.
- See TableSelectFunction for the full syntax definition.

20.4.5.3 Update Tables

You can use the `update()` method to modify one or more records in a table. The `update()` method works by filtering a query to include only the records to be updated and then applying the operations you specify to those records.

To replace a city name in the city table, pass to the `set()` method the new city name. Then, pass to the `where()` method the city name to locate and replace. The following example replaces the city Peking with Beijing.

```
mysql-js> db.city.update().set("Name", "Beijing").where("Name = 'Peking'")
Query OK, 1 item affected (0.04 sec)
```

Use the `select()` method to verify the change.

```
mysql-js> db.city.select(["ID", "Name", "CountryCode", "District", "Info"]).where("Name = 'Beijing'")
+------+-----------+-------------+----------+----------------------------+
| ID   | Name      | CountryCode | District | Info                       |
+------+-----------+-------------+----------+----------------------------+
| 1891 | Beijing   | CHN         | Peking   | {"Population": 7472000}    |
+------+-----------+-------------+----------+----------------------------+
```

```
1 row in set (0.00 sec)
```

Related Information

- See TableUpdateFunction for the full syntax definition.

20.4.5.4 Delete Tables

You can use the `delete()` method to remove some or all records from a table in a database. The X DevAPI provides additional methods to use with the `delete()` method to filter and order the records to be deleted.

Delete Records Using Conditions

The example that follows passes search conditions to the `delete()` method. All records matching the condition will be deleted from the city table. In this example, one record matches the condition.

```
mysql-js> db.city.delete().where("Name = 'Olympia'")
Query OK, 1 item affected (0.01 sec)
```

Delete the First Record

To delete the first record in the city table, use the `limit()` method with a value of 1.

```
mysql-js> db.city.delete().limit(1)
Query OK, 1 item affected (0.02 sec)
```

Delete All Records in a Table

You can delete all records in a table. To do so, use the `delete()` method without specifying a search condition.

Caution

Use care when you delete records without specifying a search condition. This action will delete all records from the table.

Related Information

- See TableDeleteFunction for the full syntax definition.
- See Section 20.4.2, “Import Database Sample” for instructions to recreate the `world_x` database.

20.4.6 Documents in Tables

In MySQL, a table may contain traditional relational data, JSON values, or both. You can combine traditional data with JSON documents by storing the documents in columns having a native `JSON` data type.

Examples in this section use the city table in the `world_x` database.

city Table Description

The city table has five columns (or fields).

```
+---------------+------------+-------+-------+---------+------------------+
| Field         | Type       | Null  | Key   | Default | Extra            |
+---------------+------------+-------+-------+---------+------------------+
| ID            | int(11)    | NO    | PRI   | null    | auto_increment   |
| Name          | char(35)   | NO    |       |         |                  |
| CountryCode   | char(3)    | NO    |       |         |                  |
| District      | char(20)   | NO    |       |         |                  |
| Info          | json       | YES   |       | null    |                  |
```

```
+--------------+------------+-------+-------+---------+-----------------+
```

Insert a Record

To insert a document into the column of a table, pass to the `values()` method a well-formed JSON document in the correct order. In the following example, a document is passed as the final value to be inserted into the Info column.

```
mysql-js> db.city.insert().
values(null, "San Francisco", "USA", "California", '{"Population":830000}')
Query OK, 1 item affected (0.01 sec)
```

Select a Record

You can issue a query with a search condition that evaluates document values in the expression.

```
mysql-js> db.city.select(["ID", "Name", "CountryCode", "District", "Info"]).
  where("CountryCode = :country and Info->'$.Population' > 1000000").
  bind('country', 'USA')
+------+---------------+-------------+---------------+-----------------------------+
| ID   | Name          | CountryCode | District      | Info                        |
+------+---------------+-------------+---------------+-----------------------------+
| 3793 | New York      | USA         | New York      | {"Population": 8008278}     |
| 3794 | Los Angeles   | USA         | California    | {"Population": 3694820}     |
| 3795 | Chicago       | USA         | Illinois      | {"Population": 2896016}     |
| 3796 | Houston       | USA         | Texas         | {"Population": 1953631}     |
| 3797 | Philadelphia  | USA         | Pennsylvania  | {"Population": 1517550}     |
| 3798 | Phoenix       | USA         | Arizona       | {"Population": 1321045}     |
| 3799 | San Diego     | USA         | California    | {"Population": 1223400}     |
| 3800 | Dallas        | USA         | Texas         | {"Population": 1188580}     |
| 3801 | San Antonio   | USA         | Texas         | {"Population": 1144646}     |
+------+---------------+-------------+---------------+-----------------------------+
9 rows in set (0.01 sec)
```

Related Information

- See Working with Relational Tables and Documents for a general overview.
- See Section 11.6, “The JSON Data Type” for a detailed description of the data type.

20.5 Quick-Start Guide: MySQL Shell for Python

This quick-start guide provides instructions to begin prototyping database applications interactively with MySQL Shell. The guide includes the following topics:

- Introduction to MySQL functionality, MySQL Shell, and the `world_x` database sample.
- Operations to manage collections and documents.
- Operations to manage relational tables.
- Operations that apply to documents within tables.

Available Quick-Start Guides

- MySQL Shell for JavaScript
- MySQL for Visual Studio

Related Information

- Chapter 19, *MySQL Shell User Guide* provides more in-depth information about MySQL Shell.
- X DevAPI User Guide provides more examples of using X DevAPI.

20.5.1 Introduction

The MySQL Shell for Python quick start provides a short but comprehensive introduction to the database functionality including the new X DevAPI, which offers a modern, integrative way to work with relational and document data, without requiring SQL knowledge from application developers.

In MySQL, tables are the native data storage container type and collections are stored internally using tables.

JSON Documents and Collections

A JSON document is a data structure composed of field/value pairs stored within a collection. The values of fields often contain other documents, arrays, and lists of documents.

```
{
    "GNP": .6,
    "IndepYear": 1967,
    "Name": "Sealand",
    "_id": "SEA",
    "demographics": {
        "LifeExpectancy": 79,
        "Population": 27
    },
    "geography": {
        "Continent": "Europe",
        "Region": "British Islands",
        "SurfaceArea": 193
    },
    "government": {
        "GovernmentForm": "Monarchy",
        "HeadOfState": "Michael Bates"
    }
}
```

Relational Tables

A table in MySQL enables you to store data organized in rows and columns. The structure of a table is defined by one or more columns with user-defined names and data types. Every row stored in the table has the same structure.

```
+------+------------------+-------------+-----------------+------------------------+
| ID   | Name             | CountryCode | District        | Info                   |
+------+------------------+-------------+-----------------+------------------------+
|    1 | Kabul            | AFG         | Kabol           |{"Population": 1780000} |
|    2 | Qandahar         | AFG         | Qandahar        |{"Population": 237500}  |
|    3 | Herat            | AFG         | Herat           |{"Population": 186800}  |
|    4 | Mazar-e-Sharif   | AFG         | Balkh           |{"Population": 127800}  |
|    5 | Amsterdam        | NLD         | Noord-Holland   |{"Population": 731200}  |
|    6 | Rotterdam        | NLD         | Zuid-Holland    |{"Population": 593321}  |
+------+------------------+-------------+-----------------+------------------------+
```

Related Information

- Chapter 19, *MySQL Shell User Guide* provides a general overview.
- See X DevAPI User Guide for development reference documentation.

20.5.2 Import Database Sample

The `world_x` database sample contains one JSON collection and a set of three relational tables:

- Collection
  - countryinfo: Information about countries in the world.

- Tables
  - country: Minimal information about countries of the world.
  - city: Information about some of the cities in those countries.
  - countrylanguage: Languages spoken in each country.

Requirements

You must install MySQL Shell with the X Protocol enabled. For instructions, see Section 20.3, “Setting Up MySQL as a Document Store”.

Start the server before you load the `world_x` database for this guide.

Download world_x Database

To prepare the `world_x` database sample, follow these steps:

1. Download world_x-db.zip.
2. Extract the installation archive to a temporary location such as `/tmp/`. Unpacking the archive results in a single file named `world_x.sql`.
3. Create the schema with the following command:

```
mysqlsh -u root --sql < /tmp/world_x-db/world_x.sql
Enter password: ****
Records: 4079  Duplicates: 0  Warnings: 0

Records: 239  Duplicates: 0  Warnings: 0

Records: 239  Duplicates: 0  Warnings: 0

Records: 984  Duplicates: 0  Warnings: 0
```

   Enter your password when prompted. A non-root account can be used as long as the account has privileges to create new databases.

   Replace `/tmp/` with the path to the `world_x.sql` file on your system.

Related Information

- Section 19.2.1, “MySQL Shell Sessions” explains session types.
- See Chapter 2, *Installing and Upgrading MySQL* for general installation assistance.

20.5.3 MySQL Shell

MySQL Shell is a unified scripting interface to MySQL Server. It supports scripting in JavaScript and Python. JavaScript is the default processing mode. In most cases, you need an account to connect to the local MySQL server instance.

Start MySQL Shell

After you have installed and started MySQL server, connect MySQL Shell to the server instance. By default, MySQL Shell connects using the X Protocol.

On the same system where the server instance is running, open a terminal window and start MySQL Shell with the following command:

```
mysqlsh name@localhost/world_x --py
```

```
Creating an Session to name@localhost/world_x
Enter password: ****
```

You may need to specify the path as appropriate.

In addition:

- *name* represents the user name of your MySQL account.
- MySQL Shell prompts you for your password.
- The `--py` option starts MySQL Shell in Python mode. If you omit `--py`, MySQL Shell starts in JavaScript mode.
- The default schema for this session is the `world_x` database. For instructions on setting up the `world_x` database sample, see Section 20.5.2, “Import Database Sample”.

The `mysql-py>` prompt indicates that the active language for this session is Python.

```
mysql-py>
```

When you run `mysqlsh` without the host argument, MySQL Shell attempts to connect to the server instance running on the localhost interface on port 33060. To specify a different host or port number, as well as other options, see the option descriptions at `mysqlsh` — The MySQL Shell.

MySQL Shell supports input-line editing as follows:

- **left-arrow** and **right-arrow** keys move horizontally within the current input line.
- **up-arrow** and **down-arrow** keys move up and down through the set of previously entered lines.
- **Backspace** deletes the character before the cursor and typing new characters enters them at the cursor position.
- **Enter** enters the current input line.

Get Help for MySQL Shell

Type `mysqlsh --help` at the prompt of your command interpreter for a list of command-line options.

```
mysqlsh --help
```

Type `\help` at the MySQL Shell prompt for a list of available commands and their descriptions.

```
mysql-py> \help
```

Type `\help` followed by a command name for detailed help about an individual MySQL Shell command. For example, to view help on the `\connect` command, type:

```
mysql-py> \help \connect
```

Quit MySQL Shell

To quit MySQL Shell, type the following command:

```
mysql-py> \quit
```

Related Information

- See Section 19.3.1, “Interactive Code Execution” for an explanation of how interactive code execution works in MySQL Shell.

- See Section 19.2, “Getting Started with MySQL Shell” to learn about session and connection alternatives.

20.5.4 Documents and Collections

In MySQL, collections contain JSON documents that you can add, find, update, and remove. Collections are containers within a schema that you create, list, and drop.

The examples in this section use the countryinfo collection in the `world_x` database. For instructions on setting up the `world_x` database sample, see Section 20.5.2, “Import Database Sample”.

Documents

In MySQL, documents are represented as JSON objects. Internally, they are stored in an efficient binary format that enables fast lookups and updates.

- Simple document format for Python:

```
{"field1": "value", "field2" : 10, "field 3": null}
```

An array of documents consists of a set of documents separated by commas and enclosed within `[` and `]` characters.

- Simple array of documents for Python:

```
[{"Name": "Aruba", "_id": "ABW"}, {"Name": "Angola", "_id": "AGO"}]
```

MySQL supports the following Python value types in JSON documents:

- numbers (integer and floating point)
- strings
- boolean (False and True)
- None
- arrays of more JSON values
- nested (or embedded) objects of more JSON values

Collections

Collections are containers for documents that share a purpose and possibly share one or more indexes. Each collection has a unique name and exists within a single schema.

The term schema is equivalent to a database, which means a group of database objects (as opposed to relational schema used to enforce structure and constraints over data). A schema does not enforce conformity on the documents in a collection.

In this quick-start guide:

- Basic objects include:

Object form	Description
`db`	`db` is a global variable assigned to the current active schema that you specified on the command line. You can type `db` in MySQL Shell to print a description of the object, which in this case will be the name of the schema it represents.

Object form	Description
`db.get_collections()`	db.get_collections() holds a list of collections in the schema. Use the list to get references to collection objects, iterate over them, and so on.

- Basic operations scoped by collections include:

Operation form	Description
`db.name.add()`	The add() method inserts one document or a list of documents into the named collection.
`db.name.find()`	The find() method returns some or all documents in the named collection.
`db.name.modify()`	The modify() method updates documents in the named collection.
`db.name.remove()`	The remove() method deletes one document or a list of documents from the named collection.

Related Information

- See Working with Collections for a general overview.
- CRUD EBNF Definitions provides a complete list of operations.

20.5.4.1 Create, List, and Drop Collections

In MySQL Shell, you can create new collections, get a list of the existing collections in a schema, and remove an existing collection from a schema. Collection names are case-sensitive and each collection name must be unique.

Confirm the Schema

To show the value that is assigned to the schema variable, type `db`.

```
mysql-py> db
<Schema:world_x>
```

If the schema value is not `Schema:world_x`, then set the `db` variable as follows:

```
mysql-py> \use world_x
Schema `world_x` accessible through db.
```

Create a Collection

To create a new collection in an existing schema, use the `create_collection()` method.

The following example creates in the `world_x` database a collection called flags. The method returns a collection object.

```
mysql-py> db.create_collection("flags")
<Collection:flags>
```

List Collections

To display all collections in the `world_x` database, use the `get_collections()` method on the schema object. Collections returned by the server appear between brackets.

```
mysql-py> db.get_collections()
[
```

```
    <Collection:countryinfo>,
    <Collection:flags>
]
```

Drop a Collection

To drop an existing collection from a database, use the `drop_collection()` method on the session object. For example, to drop the flags collection from the `world_x` database, type:

```
mysql-py> session.drop_collection("world_x", "flags")
Query OK (0.04 sec)
```

Related Information

- See Section 19.2.2.6, “Connections in JavaScript and Python” to learn more about the session object.
- See Collection Objects for more examples.

20.5.4.2 Add Documents

You can use the `add()` method to insert one document or a list documents into an existing collection using MySQL Shell. All examples in this section use the countryinfo collection.

Confirm the Schema

To show the value that is assigned to the schema variable, type`db`.

```
mysql-py> db
<Schema:world_x>
```

If the schema value is not `Schema:world_x`, then set the `db` variable as follows:

```
mysql-py> \use world_x
Schema `world_x` accessible through db.
```

Add a Document

Insert the following document into the countryinfo collection. Press **Enter** twice to insert the document.

```
mysql-py> db.countryinfo.add(
 {
    "GNP": .6,
    "IndepYear": 1967,
    "Name": "Sealand",
    "_id": "SEA",
    "demographics": {
        "LifeExpectancy": 79,
        "Population": 27
    },
    "geography": {
        "Continent": "Europe",
        "Region": "British Islands",
        "SurfaceArea": 193
    },
    "government": {
        "GovernmentForm": "Monarchy",
        "HeadOfState": "Michael Bates"
    }
  }
)
Query OK, 1 item affected (0.02 sec)
```

The method returns the status of the operation.

Each document requires an identifier field called `_id`. The value of the `_id` field must be unique among all documents in the same collection. If the document passed to the `add()` method does not contain the `_id` field, MySQL Shell automatically inserts a field into the document and sets the value to a generated universal unique identifier (UUID).

Related Information

- See CollectionAddFunction for the full syntax definition.

20.5.4.3 Find Documents

You can use the `find()` method to query for and return documents from a collection in a database. MySQL Shell provides additional methods to use with the `find()` method to filter and sort the returned documents.

MySQL provides the following operators to specify search conditions: `OR` (`||`), `AND` (`&&`), `XOR`, `IS`, `NOT`, `BETWEEN`, `IN`, `LIKE`, `!=`, `<>`, `>`, `>=`, `<`, `<=`, `&`, `|`, `<<`, `>>`, `+`, `-`, `*`, `/`, `~`, and `%`.

Find All Documents in a Collection

To return all documents in a collection, use the `find()` method without specifying search conditions. For example, the following operation returns all documents in the countryinfo collection.

```
mysql-py> db.countryinfo.find()
[
     {
          "GNP": 828,
          "IndepYear": null,
          "Name": "Aruba",
          "_id": "ABW",
          "demographics": {
              "LifeExpectancy": 78.4000015258789,
              "Population": 103000
          },
          "geography": {
              "Continent": "North America",
              "Region": "Caribbean",
              "SurfaceArea": 193
          },
          "government": {
              "GovernmentForm": "Nonmetropolitan Territory of The Netherlands",
              "HeadOfState": "Beatrix"
          }
          ...
      }
 ]
240 documents in set (0.00 sec)
```

The method produces results that contain operational information in addition to all documents in the collection.

An empty set (no matching documents) returns the following information:

```
Empty set (0.00 sec)
```

Filter Searches

You can include search conditions with the `find()` method. The syntax for expressions that form a search condition is the same as that of traditional MySQL. You must enclose all expressions in quotes.

All examples in this section use the countryinfo collection in the `world_x` database. For the sake of brevity, some of the examples do not display output.

A simple search condition consists of the `_id` field and unique identifier of a document. The following example returns a single document matching the identifier string:

```
mysql-py> db.countryinfo.find("_id = 'AUS'")
[
    {
        "GNP": 351182,
        "IndepYear": 1901,
        "Name": "Australia",
        "_id": "AUS",
        "demographics": {
            "LifeExpectancy": 79.80000305175781,
            "Population": 18886000
        },
        "geography": {
            "Continent": "Oceania",
            "Region": "Australia and New Zealand",
            "SurfaceArea": 7741220
        },
        "government": {
            "GovernmentForm": "Constitutional Monarchy, Federation",
            "HeadOfState": "Elisabeth II"
        }
    }
]
1 document in set (0.01 sec)
```

The following example searches for all countries that have a GNP higher than $500 billion. The countryinfo collection measures GNP in units of million.

```
mysql-py> db.countryinfo.find("GNP > 500000")
...[output removed]
10 documents in set (0.00 sec)
```

The Population field in the following query is embedded within the demographics object. To access the embedded field, use a period between demographics and Population to identify the relationship. Document and field names are case sensitive.

```
mysql-py> db.countryinfo.find("GNP > 500000 and demographics.Population < 100000000")
...[output removed]
6 documents in set (0.00 sec)
```

Arithmetic operators in the following expression are used to query for countries with a GNP per capita higher than $30000. Search conditions can include arithmetic operators and most MySQL functions.

Note

Seven documents in the countryinfo collection have a population value of zero. Warning messages appear at the end of the output.

```
mysql-py> db.countryinfo.find("GNP*1000000/demographics.Population > 30000")
...[output removed]
9 documents in set, 7 warnings (0.00 sec)
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
Warning (Code 1365): Division by 0
```

You can separate a value from the search condition by using the `bind()` method. For example, instead of specifying a hard-coded country name as the condition, substitute a named placeholder consisting of a colon followed by a name that begins with a letter, such as *country*. Then include the placeholder and value in the `bind()` method as follows:

```
mysql-py> db.countryinfo.find("Name = :country").bind("country", "Italy")
```

```
[
    {
        "GNP": 1161755,
        "IndepYear": 1861,
        "Name": "Italy",
        "_id": "ITA",
        "demographics": {
            "LifeExpectancy": 79,
            "Population": 57680000
        },
        "geography": {
            "Continent": "Europe",
            "Region": "Southern Europe",
            "SurfaceArea": 301316
        },
        "government": {
            "GovernmentForm": "Republic",
            "HeadOfState": "Carlo Azeglio Ciampi"
        }
    }
]
1 document in set (0.01 sec)
```

Tip

Within a program, binding enables you to specify placeholders in your expressions, which are filled in with values before execution and can benefit from automatic escaping, as appropriate.

Always use binding to sanitize input. Avoid introducing values in queries using string concatenation, which can produce invalid input and, in some cases, can cause security issues.

Project Results

You can return specific fields of a document, instead of returning all the fields. The following example returns the GNP and Name fields of all documents in the countryinfo collection matching the search conditions.

Use the `fields()` method to pass the list of fields to return.

```
mysql-py> db.countryinfo.find("GNP > 5000000").fields(["GNP", "Name"])
[
    {
        "GNP": 8510700,
        "Name": "United States"
    }
]
1 document in set (0.00 sec)
```

In addition, you can alter the returned documents—adding, renaming, nesting and even computing new field values—with an expression that describes the document to return. For example, alter the names of the fields with the following expression to return only two documents.

```
mysql-py> db.countryinfo.find().\
fields(mysqlx.expr('{"Name": upper(Name), "GNPPerCapita": GNP*1000000/demographics.Population}')).\
limit(2)
[
    {
        "GNPPerCapita": 8038.834951456311,
        "Name": "ARUBA"
    },
    {
        "GNPPerCapita": 263.0281690140845,
        "Name": "AFGHANISTAN"
    }
```

```
]
2 documents in set (0.00 sec)
```

Limit, Sort, and Skip Results

You can apply the `limit()`, `sort()`, and `skip()` methods to manage the number and order of documents returned by the `find()` method.

To specify the number of documents included in a result set, append the `limit()` method with a value to the `find()` method. The following query returns the first five documents in the countryinfo collection.

```
mysql-py> db.countryinfo.find().limit(5)
... [output removed]
5 documents in set (0.00 sec)
```

To specify an order for the results, append the `sort()` method to the `find()` method. Pass to the `sort()` method a list of one or more fields to sort by and, optionally, the descending (`desc`) or ascending (`asc`) attribute as appropriate. Ascending order is the default order type.

For example, the following query sorts all documents by the IndepYear field and then returns the first eight documents in descending order.

```
mysql-py> db.countryinfo.find().sort(["IndepYear desc"]).limit(8)
... [output removed]
8 documents in set (0.00 sec)
```

By default, the `limit()` method starts from the first document in the collection. You can use the `skip()` method to change the starting document. For example, to ignore the first document and return the next eight documents matching the condition, pass to the `skip()` method a value of 1.

```
mysql-py> db.countryinfo.find().sort(["IndepYear desc"]).limit(8).skip(1)
... [output removed]
8 documents in set (0.00 sec)
```

Related Information

- The MySQL Reference Manual provides detailed documentation on functions and operators.
- See CollectionFindFunction for the full syntax definition.

20.5.4.4 Modify Documents

You can use the `modify()` method to update one or more documents in a collection. The X DevAPI provides additional methods for use with the `modify()` method to:

- Set and unset fields within documents.
- Append, insert, and delete arrays.
- Bind, limit, and sort the documents to be modified.

Set and Unset Fields

The `modify()` method works by filtering a collection to include only the documents to be modified and then applying the operations that you specify to those documents.

In the following example, the `modify()` method uses the search condition to identify the document to change and then the `set()` method replaces two values within the nested demographics object.

```
mysql-py> db.countryinfo.modify("_id = 'SEA'").\
```

```
set("demographics", {"LifeExpectancy": 78, "Population": 28})
Query OK, 1 item affected (0.04 sec)
```

After you modify a document, use the `find()` method to verify the change.

To remove content from a document, use the `modify()` and `unset()` methods. For example, the following query removes the GNP from a document that matches the search condition.

```
mysql-py> db.countryinfo.modify("Name = 'Sealand'").unset("GNP")
Query OK, 1 item affected (0.01 sec)
```

Use the `find()` method to verify the change.

```
mysql-py> db.countryinfo.find("Name = 'Sealand'")
[
    {
        "IndepYear": 1967,
        "Name": "Sealand",
        "_id": "SEA",
        "demographics": {
            "LifeExpectancy": 78,
            "Population": 28
        },
        "geography": {
            "Continent": "Europe",
            "Region": "British Islands",
            "SurfaceArea": 193
        },
        "government": {
            "GovernmentForm": "Monarchy",
            "HeadOfState": "Michael Bates"
        }
    }
]
1 document in set (0.00 sec)
```

Append, Insert, and Delete Arrays

To append an element to an array field, or insert, or delete elements in an array, use the `array_append()`, `array_insert()`, or `array_delete()` methods. The following examples modify the countryinfo collection to enable tracking of international airports.

The first example uses the `modify()` and `set()` methods to create a new Airports field in all documents.

Caution

Use care when you modify documents without specifying a search condition. This action will modify all documents in the collection.

```
mysql-py> db.countryinfo.modify().set("Airports", [])
Query OK, 240 items affected (0.07 sec)
```

With the Airports field added, the next example uses the `array_append()` method to add a new airport to one of the documents. *$.Airports* in the following example represents the Airports field of the current document.

```
mysql-py> db.countryinfo.modify("Name = 'France'").array_append("$.Airports", "ORY")
Query OK, 1 item affected (0.02 sec)
```

Use **`db.countryinfo.find("Name = 'France'")`** to see the change.

To insert an element at a different position in the array, use the `array_insert()` method to specify which index to insert in the path expression. In this case, the index is 0, or the first element in the array.

```
mysql-py> db.countryinfo.modify("Name = 'France'").array_insert("$.Airports[0]", "CDG")
Query OK, 1 item affected (0.04 sec)
```

To delete an element from the array, you must pass to the `array_delete()` method the index of the element to be deleted.

```
mysql-py> db.countryinfo.modify("Name = 'France'").array_delete("$.Airports[1]")
Query OK, 1 item affected (0.03 sec)
```

Related Information

- The MySQL Reference Manual provides instructions to help you search for and modify JSON values.
- See CollectionModifyFunction for the full syntax definition.

20.5.4.5 Remove Documents

You can use the `remove()` method to delete some or all documents from a collection in a database. The X DevAPI provides additional methods for use with the `remove()` method to filter and sort the documents to be removed.

Remove Documents Using Conditions

The example that follows passes a search condition to the `remove()` method. All documents matching the condition will be removed from the countryinfo collection. In this example, one document matches the condition.

```
mysql-py> db.countryinfo.remove("_id = 'SEA'")
Query OK, 1 item affected (0.02 sec)
```

Remove the First Document

To remove the first document in the countryinfo collection, use the `limit()` method with a value of 1.

```
mysql-py> db.countryinfo.remove().limit(1)
Query OK, 1 item affected (0.03 sec)
```

Remove the Last Document in an Order

The following example removes the last document in the countryinfo collection by country name.

```
mysql-py> db.countryinfo.remove().sort(["Name desc"]).limit(1)
Query OK, 1 item affected (0.02 sec)
```

Remove All Documents in a Collection

You can remove all documents in a collection. To do so, use the `remove()` method without specifying a search condition.

Caution

Use care when you remove documents without specifying a search condition. This action will delete all documents from the collection.

Related Information

- See CollectionRemoveFunction for the full syntax definition.
- See Section 20.5.2, “Import Database Sample” for instructions to recreate the `world_x` database.

20.5.4.6 Create and Drop Indexes

Indexes are used to find documents with specific field values quickly. Without an index, MySQL must begin with the first document and then read through the entire collection to find the relevant fields. The larger the collection, the more this costs. If a collection is large and queries on a specific field are common, then consider creating an index on a specific field inside a document.

For example, the following query will perform better with an index:

```
mysql-js> db.countryinfo.find("demographics.Population < 100")
...[output removed]
8 documents in set (0.00 sec)
```

The `create_index()` method creates an index that you can define as nonunique or unique. Use the `field()` method to chain the fields that should be indexed. The `execute()` method is required to create or drop an index.

In MySQL, the `_id` field is equivalent to a primary key by default.

Add a Nonunique Index

To create a nonunique index, pass to the `create_index()` method an index name. Duplicate index names are prohibited.

In the following example, the first parameter of the `field()` method specifies the Population field inside the demographics object and the next parameter indicates that the field should be indexed as an Integer numeric value. The last parameter indicates whether the field should require the NOT NULL constraint. If the value is `False`, the field can contain `NULL` values.

```
mysql-js> db.countryinfo.create_index("pop").\
field("demographics.Population", "INTEGER", False).execute()
Query OK (0.04 sec)
```

Add a Unique Index

To create a unique index, pass to the `create_index()` method an index name and the `mysqlx.IndexType.UNIQUE` type. Country `"Name"` is another common field in the countryinfo collection to index. In the following example, `"Text(40)"` represents the number of characters to index and `True` indicates that the field cannot contain any `NULL` values.

```
mysql-js> db.countryinfo.create_index("name", mysqlx.IndexType.UNIQUE).\
field("Name", "TEXT(40)", True).execute()
Query OK (0.04 sec)
```

Drop an Index

To drop an index, pass to the `drop_index()` method the name of the index to drop. For example, you can drop the “pop” index as follows:

```
mysql-js> db.countryinfo.drop_index("pop").execute()
Query OK (0.58 sec)
```

Related Information

- See Collection Index Management Functions for the full syntax definition.

20.5.5 Relational Tables

You can use MySQL Shell to manipulate not just JSON documents, but also relational tables.

In MySQL, each relational table is associated with a particular storage engine. The examples in this section use `InnoDB` tables in the `world_x` database.

Confirm the Schema

To show the value that is assigned to the schema variable, type `db`.

```
mysql-py> db
<Schema:world_x>
```

If the schema value is not the `Schema:world_x` database, then set the `db` variable as follows:

```
mysql-py> \use world_x
Schema `world_x` accessible through db.
```

Show All Tables

To display all relational tables in the `world_x` database, use the `get_tables()` method on the schema object.

```
mysql-py> db.get_tables()
{
    "city": <Table:city>,
    "country": <Table:country>,
    "countrylanguage": <Table:countrylanguage>
}
```

Basic Table Operations

Basic operations scoped by tables include:

Operation form	Description
`db.name.insert()`	The insert() method inserts one or more records into the named table.
`db.name.select()`	The select() method returns some or all records in the named table.
`db.name.update()`	The update() method updates records in the named table.
`db.name.delete()`	The delete() method deletes one or more records from the named table.

Related Information

- See Working with Relational Tables for a general overview.
- CRUD EBNF Definitions provides a complete list of operations.
- See Section 20.5.2, “Import Database Sample” for instructions on setting up the `world_x` database sample.

20.5.5.1 Insert Records into Tables

You can use the `insert()` method with the `values()` method to insert records into an existing relational table. The `insert()` method accepts individual columns or all columns in the table. Use one or more `values()` methods to specify the values to be inserted.

Insert a Complete Record

To insert a complete record, pass to the `insert()` method all columns in the table. Then pass to the `values()` method one value for each column. For example, to add a new record to the city table in the `world_x` database, insert the following record and press **Enter** twice.

```
mysql-py> db.city.insert("ID", "Name", "CountryCode", "District", "Info").\
values(None, "Olympia", "USA", "Washington", '{"Population": 5000}')
```

```
Query OK, 1 item affected (0.01 sec)
```

The city table has five columns: ID, Name, CountryCode, District, and Info. Each value must match the data type of the column it represents.

Insert a Partial Record

The following example inserts values into the ID, Name, and CountryCode columns of the city table.

```
mysql-py> db.city.insert("ID", "Name", "CountryCode").\
values(None, "Little Falls", "USA").values(None, "Happy Valley", "USA")
Query OK, 2 item affected (0.03 sec)
```

When you specify columns using the `insert()` method, the number of values must match the number of columns. In the previous example, you must supply three values to match the three columns specified.

Related Information

- See TableInsertFunction for the full syntax definition.

20.5.5.2 Select Tables

You can use the `select()` method to query for and return records from a table in a database. The X DevAPI provides additional methods to use with the `select()` method to filter and sort the returned records.

MySQL provides the following operators to specify search conditions: `OR` (`||`), `AND` (`&&`), `XOR`, `IS`, `NOT`, `BETWEEN`, `IN`, `LIKE`, `!=`, `<>`, `>`, `>=`, `<`, `<=`, `&`, `|`, `<<`, `>>`, `+`, `-`, `*`, `/`, `~`, and `%`.

Select All Records

To issue a query that returns all records from an existing table, use the `select()` method without specifying search conditions. The following example selects all records from the city table in the `world_x` database.

Note

Limit the use of the empty `select()` method to interactive statements. Always use explicit column-name selections in your application code.

```
mysql-py> db.city.select()
+------+------------+-------------+------------+-------------------------+
| ID   | Name       | CountryCode | District   | Info                    |
+------+------------+-------------+------------+-------------------------+
|    1 | Kabul      | AFG         | Kabol      |{"Population": 1780000}  |
|    2 | Qandahar   | AFG         | Qandahar   |{"Population": 237500}   |
|    3 | Herat      | AFG         | Herat      |{"Population": 186800}   |
...    ...          ...           ...          ...
| 4079 | Rafah      | PSE         | Rafah      |{"Population": 92020}    |
+------+------- ----+-------------+------------+-------------------------+
4082 rows in set (0.01 sec)
```

An empty set (no matching records) returns the following information:

```
Empty set (0.00 sec)
```

Filter Searches

To issue a query that returns a set of table columns, use the `select()` method and specify the columns to return between square brackets. This query returns the Name and CountryCode columns from the city table.

```
mysql-py> db.city.select(["Name", "CountryCode"])
+-------------------+-------------+
| Name              | CountryCode |
+-------------------+-------------+
| Kabul             | AFG         |
| Qandahar          | AFG         |
| Herat             | AFG         |
| Mazar-e-Sharif    | AFG         |
| Amsterdam         | NLD         |
...                 ...
| Rafah             | PSE         |
| Olympia           | USA         |
| Little Falls      | USA         |
| Happy Valley      | USA         |
+-------------------+-------------+
4082 rows in set (0.00 sec)
```

To issue a query that returns rows matching specific search conditions, use the `where()` method to include those conditions. For example, the following example returns the names and country codes of the cities that start with the letter Z.

```
mysql-py> db.city.select(["Name", "CountryCode"]).where("Name like 'Z%'")
+-------------------+-------------+
| Name              | CountryCode |
+-------------------+-------------+
| Zaanstad          | NLD         |
| Zoetermeer        | NLD         |
| Zwolle            | NLD         |
| Zenica            | BIH         |
| Zagazig           | EGY         |
| Zaragoza          | ESP         |
| Zamboanga         | PHL         |
| Zahedan           | IRN         |
| Zanjan            | IRN         |
| Zabol             | IRN         |
| Zama              | JPN         |
| Zhezqazghan       | KAZ         |
| Zhengzhou         | CHN         |
...                 ...
| Zeleznogorsk      | RUS         |
+-------------------+-------------+
59 rows in set (0.00 sec)
```

You can separate a value from the search condition by using the `bind()` method. For example, instead of using "Name = 'Z%' " as the condition, substitute a named placeholder consisting of a colon followed by a name that begins with a letter, such as *name*. Then include the placeholder and value in the `bind()` method as follows:

```
mysql-py> db.city.select(["Name", "CountryCode"]).\
              where("Name like :name").bind("name", "Z%")
```

Tip

Within a program, binding enables you to specify placeholders in your expressions, which are filled in with values before execution and can benefit from automatic escaping, as appropriate.

Always use binding to sanitize input. Avoid introducing values in queries using string concatenation, which can produce invalid input and, in some cases, can cause security issues.

Project Results

To issue a query using the `AND` operator, add the operator between search conditions in the `where()` method.

```
mysql-py> db.city.select(["Name", "CountryCode"]).\
              where("Name like 'Z%' and CountryCode = 'CHN'")
+---------------+-------------+
| Name          | CountryCode |
+---------------+-------------+
| Zhengzhou     | CHN         |
| Zibo          | CHN         |
| Zhangjiakou   | CHN         |
| Zhuzhou       | CHN         |
| Zhangjiang    | CHN         |
| Zigong        | CHN         |
| Zaozhuang     | CHN         |
...             ...
| Zhangjiagang  | CHN         |
+---------------+-------------+
22 rows in set (0.01 sec)
```

To specify multiple conditional operators, you can enclose the search conditions in parenthesis to change the operator precedence. The following example demonstrates the placement of `AND` and `OR` operators.

```
mysql-py> db.city.select(["Name", "CountryCode"]).\
where("Name like 'Z%' and (CountryCode = 'CHN' or CountryCode = 'RUS')")
+------------------+-------------+
| Name             | CountryCode |
+------------------+-------------+
| Zhengzhou        | CHN         |
| Zibo             | CHN         |
| Zhangjiakou      | CHN         |
| Zhuzhou          | CHN         |
...                ...
| Zeleznogorsk     | RUS         |
+------------------+-------------+
29 rows in set (0.01 sec)
```

Limit, Order, and Offset Results

You can apply the `limit()`, `order_by()`, and `offset()` methods to manage the number and order of records returned by the `select()` method.

To specify the number of records included in a result set, append the `limit()` method with a value to the `select()` method. For example, the following query returns the first five records in the country table.

```
mysql-py> db.country.select(["Code", "Name"]).limit(5)
+------+-------------+
| Code | Name        |
+------+-------------+
| ABW  | Aruba       |
| AFG  | Afghanistan |
| AGO  | Angola      |
| AIA  | Anguilla    |
| ALB  | Albania     |
+------+-------------+
5 rows in set (0.00 sec)
```

To specify an order for the results, append the `order_by()` method to the `select()` method. Pass to the `order_by()` method a list of one or more columns to sort by and, optionally, the descending (`desc`) or ascending (`asc`) attribute as appropriate. Ascending order is the default order type.

For example, the following query sorts all records by the Name column and then returns the first three records in descending order .

```
mysql-py> db.country.select(["Code", "Name"]).order_by(["Name desc"]).limit(3)
+------+------------+
| Code | Name       |
+------+------------+
| ZWE  | Zimbabwe   |
| ZMB  | Zambia     |
| YUG  | Yugoslavia |
+------+------------+
3 rows in set (0.00 sec)
```

By default, the `limit()` method starts from the first record in the table. You can use the `offset()` method to change the starting record. For example, to ignore the first record and return the next three records matching the condition, pass to the `offset()` method a value of 1.

```
mysql-py> db.country.select(["Code", "Name"]).order_by(["Name desc"]).limit(3).offset(1)
+------+------------+
| Code | Name       |
+------+------------+
| ZMB  | Zambia     |
| YUG  | Yugoslavia |
| YEM  | Yemen      |
+------+------------+
3 rows in set (0.00 sec)
```

Related Information

- The MySQL Reference Manual provides detailed documentation on functions and operators.
- See TableSelectFunction for the full syntax definition.

20.5.5.3 Update Tables

You can use the `update()` method to modify one or more records in a table. The `update()` method works by filtering a query to include only the records to be updated and then applying the operations you specify to those records.

To replace a city name in the city table, pass to the `set()` method the new city name. Then, pass to the `where()` method the city name to locate and replace. The following example replaces the city Peking with Beijing.

```
mysql-py> db.city.update().set("Name", "Beijing").where("Name = 'Peking'")
Query OK, 1 item affected (0.04 sec)
```

Use the `select()` method to verify the change.

```
mysql-py> db.city.select(["ID", "Name", "CountryCode", "District", "Info"]).where("Name = 'Beijing'")
+------+-----------+-------------+----------+-----------------------------+
| ID   | Name      | CountryCode | District | Info                        |
+------+-----------+-------------+----------+-----------------------------+
| 1891 | Beijing   | CHN         | Peking   | {"Population": 7472000}     |
+------+-----------+-------------+----------+-----------------------------+
1 row in set (0.00 sec)
```

Related Information

- See TableUpdateFunction for the full syntax definition.

20.5.5.4 Delete Tables

You can use the `delete()` method to remove some or all records from a table in a database. The X DevAPI provides additional methods to use with the `delete()` method to filter and order the records to be deleted.

Delete Records Using Conditions

The example that follows passes search conditions to the `delete()` method. All records matching the condition will be deleted from the city table. In this example, one record matches the condition.

```
mysql-py> db.city.delete().where("Name = 'Olympia'")
Query OK, 1 item affected (0.01 sec)
```

Delete the First Record

To delete the first record in the city table, use the `limit()` method with a value of 1.

```
mysql-py> db.city.delete().limit(1)
Query OK, 1 item affected (0.02 sec)
```

Delete All Records in a Table

You can delete all records in a table. To do so, use the `delete()` method without specifying a search condition.

Caution

Use care when you delete records without specifying a search condition. This action will delete all records from the table.

Related Information

- See TableDeleteFunction for the full syntax definition.
- See Section 20.5.2, “Import Database Sample” for instructions to recreate the `world_x` database.

20.5.6 Documents in Tables

In MySQL, a table may contain traditional relational data, JSON values, or both. You can combine traditional data with JSON documents by storing the documents in columns having a native `JSON` data type.

Examples in this section use the city table in the `world_x` database.

city Table Description

The city table has five columns (or fields).

```
+--------------+------------+-------+-------+---------+------------------+
| Field        | Type       | Null  | Key   | Default | Extra            |
+--------------+------------+-------+-------+---------+------------------+
| ID           | int(11)    | NO    | PRI   | null    | auto_increment   |
| Name         | char(35)   | NO    |       |         |                  |
| CountryCode  | char(3)    | NO    |       |         |                  |
| District     | char(20)   | NO    |       |         |                  |
| Info         | json       | YES   |       | null    |                  |
+--------------+------------+-------+-------+---------+------------------+
```

Insert a Record

To insert a document into the column of a table, pass to the `values()` method a well-formed JSON document in the correct order. In the following example, a document is passed as the final value to be inserted into the Info column.

```
 mysql-py> db.city.insert().\
values(None, "San Francisco", "USA", "California", '{"Population":830000}')
```

```
Query OK, 1 item affected (0.01 sec)
```

Select a Record

You can issue a query with a search condition that evaluates document values in the expression.

```
mysql-py> db.city.select(["ID", "Name", "CountryCode", "District", "Info"]).\
  where("CountryCode = :country and Info->'$.Population' > 1000000").\
  bind('country', 'USA')
+------+--------------+-------------+----------------+-----------------------------+
| ID   | Name         | CountryCode | District       | Info                        |
+------+--------------+-------------+----------------+-----------------------------+
| 3793 | New York     | USA         | New York       | {"Population": 8008278}     |
| 3794 | Los Angeles  | USA         | California     | {"Population": 3694820}     |
| 3795 | Chicago      | USA         | Illinois       | {"Population": 2896016}     |
| 3796 | Houston      | USA         | Texas          | {"Population": 1953631}     |
| 3797 | Philadelphia | USA         | Pennsylvania   | {"Population": 1517550}     |
| 3798 | Phoenix      | USA         | Arizona        | {"Population": 1321045}     |
| 3799 | San Diego    | USA         | California     | {"Population": 1223400}     |
| 3800 | Dallas       | USA         | Texas          | {"Population": 1188580}     |
| 3801 | San Antonio  | USA         | Texas          | {"Population": 1144646}     |
+------+--------------+-------------+----------------+-----------------------------+
9 rows in set (0.01 sec)
```

Related Information

- See Working with Relational Tables and Documents for a general overview.
- See Section 11.6, “The JSON Data Type” for a detailed description of the data type.

20.6 Quick-Start Guide: MySQL for Visual Studio

This section explains how to use MySQL Shell to script a server using MySQL for Visual Studio.

Introduction

MySQL for Visual Studio provides access to MySQL objects and data without forcing developers to leave Visual Studio. Designed and developed as a Visual Studio package, MySQL for Visual Studio integrates directly into Server Explorer providing a seamless experience for setting up new connections and working with database objects.

The following MySQL for Visual Studio features are available as of version 2.0.2:

- JavaScript and Python code editors, where scripts in those languages can be executed to query data from a MySQL database.
- Better integration with the Server Explorer to open MySQL, JavaScript, and Python code editors directly from a connected MySQL instance.
- A newer user interface for displaying query results, where different views are presented from result sets returned by a MySQL Server like:
  - Multiple tabs for each result set returned by an executed query.
  - Results view, where the information can be seen in grid, tree, or text representation for JSON results.
  - Field types view, where information about the columns of a result set is shown, such as names, data types, character sets, and more.
  - Query statistics view, displaying information about the executed query such as execution times, processed rows, index and temporary tables usage, and more.

- Execution plan view, displaying an explanation of the query execution done internally by the MySQL Server.

Getting Started

The requirements are MySQL for Visual Studio 2.0.2 or higher, and Visual Studio 2010 or higher. X DevAPI support requires MySQL Server 5.7.12 or higher with the X plugin enabled.

Opening a Code Editor

Before opening a code editor that can execute queries against a MySQL server, a connection needs to be established:

1. Open the Server Explorer pane through the View menu, or with **Control** + **W**, **K**.
2. Right-click on the Data Connections node, select **Add Connection...**.
3. In the **Add Connection** dialog, make sure the MySQL Data Provider is being used and fill in all the information.

Note

To enter the port number, click **Advanced...** and set the Port among the list of connection properties.

4. Click **Test Connection** to ensure you have a valid connection, then click **OK**.
5. Right-click your newly created connection, select **New MySQL Script** and then the language for the code editor you want to open.

For existing MySQL connections, to create a new editor you need only to do the last step.

Using the Code Editor

The MySQL script editors have a toolbar at the top where information about the session is displayed, along with the actions that can be executed.

Note

Note the first two buttons in the toolbar represent a way to connect or disconnect from a MySQL server. If the editor was opened from the Server Explorer, the connection will be already established for the new editor window.

The third button is the **Run** button, the script contained in the editor window is executed by clicking it and results from the script execution are displayed below the script window.

Note

Some commands in the MySQL Shell can be executed without appending `execute()` while in interactive mode. In MySQL for Visual Studio, these commands do require `execute()`. In other words, append ".execute()" to execute commands.

20.7 X Plugin

X Plugin extends MySQL Server to be able to function as a document store. Running X Plugin enables MySQL Server to communicate with clients using the X Protocol, which is designed to expose the ACID compliant storage abilities of MySQL as a document store. Documents are stored in JSON format and

enable schema-less storage. Using the X DevAPI you can use a NoSQL-like syntax to execute Create, Read, Update, Delete (CRUD) operations against these documents.

This section explains how to configure and monitor the X Plugin in MySQL Server.

20.7.1 Using Secure Connections with X Plugin

This section explains how to configure X Plugin to use secure connections. For more background information, see Section 6.4, “Using Secure Connections”.

X Plugin has its own SSL settings which can differ from those used with MySQL Server. This means that X Plugin can be configured with a different SSL key, certificate, and certificate authorities file than MySQL Server. Similarly, X Plugin has its own SSL status variables calculated independently from the MySQL Server SSL related variables. By default the X Plugin SSL configuration is taken from the `mysqlx_ssl_*` variables, described at Section 20.7.2.2, “X Plugin System Variables and Options”. If no configuration is provided using the `mysqlx_ssl_*` variables, X Plugin falls back to using the MySQL Server SSL system variables. This means you can choose to either have separate SSL configurations for MySQL Protocol and X Protocol connections by configuring each separately, or share the SSL configuration between MySQL Protocol and X Protocol connections by only configuring the `ssl-*` variables.

On a server with X Plugin installed, to configure MySQL Protocol and X Protocol connections with separate SSL configurations use both the `ssl-*` and `mysqlx-ssl-*` variables in `my.cnf`:

```
[mysqld]
ssl-ca=ca1.pem
ssl-cert=server-cert1.pem
ssl-key=server-key1.pem

mysqlx-ssl-ca=ca2.pem
mysqlx-ssl-cert=server-cert2.pem
mysqlx-ssl-key=server-key2.pem
```

The available `mysqlx_ssl_*` variables mirror the SSL variables in MySQL Server, so the files and techniques described for configuring MySQL Server to use SSL at Section 6.4.4, “Configuring MySQL to Use Secure Connections” are relevant to configuring X Plugin to use secure connections. An exception to this rule is the `mysqlx_ssl` variable, which defaults to `1` whenever any of the other X Protocol SSL related variables is configured. In other words, if you configure any of the `mysqlx_ssl_*` variables there is no need to configure `mysqlx_ssl`.

You can configure the TLS versions used by X Protocol SSL connections using the `tls_version` system variable. The TLS version used by MySQL Protocol and X Protocol connections is therefore the same TLS version.

Encryption per connection is optional, but a specific user can be forced to use encryption for X Protocol and MySQL Protocol connections. You configure such a user by issuing a `GRANT` statement with the `REQUIRE` option. For more details see Section 13.7.1.6, “GRANT Syntax”. Alternatively all X Protocol and MySQL Protocol connections can be forced to use encryption by setting `require_secure_transport`.

20.7.2 X Plugin Options and Variables

This section describes the command options and system variables which configure X Plugin. If values specified at startup time are incorrect, X Plugin could fail to initialize properly and the server does not load it. In this case, the server could also produce error messages for other X Plugin settings because it cannot recognize them.

20.7.2.1 X Plugin Option and Variable Reference

This table provides an overview of the command options, and system and status variables provided by X Plugin.

Table 20.1 X Plugin Option and Variable Reference

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Mysqlx_address				Yes	Global	No
mysqlx_bind_address	Yes	Yes	Yes		Global	No
Mysqlx_bytes_received				Yes	Both	No
Mysqlx_bytes_sent				Yes	Both	No
mysqlx_connect_timeout	Yes	Yes	Yes		Global	Yes
Mysqlx_connection_accept_errors				Yes	Both	No
Mysqlx_connection_errors				Yes	Both	No
Mysqlx_connections_accepted				Yes	Global	No
Mysqlx_connections_closed				Yes	Global	No
Mysqlx_connections_rejected				Yes	Global	No
Mysqlx_crud_create_view				Yes	Both	No
Mysqlx_crud_delete				Yes	Both	No
Mysqlx_crud_drop_view				Yes	Both	No
Mysqlx_crud_find				Yes	Both	No
Mysqlx_crud_insert				Yes	Both	No
Mysqlx_crud_modify_view				Yes	Both	No
Mysqlx_crud_update				Yes	Both	No
Mysqlx_errors_sent				Yes	Both	No
Mysqlx_errors_unknown_message_type				Yes	Both	No
Mysqlx_expect_close				Yes	Both	No
Mysqlx_expect_open				Yes	Both	No
mysqlx_idle_worker_thread_timeout	Yes	Yes				
Mysqlx_init_error				Yes	Both	No
mysqlx_max_allowed_packet	Yes	Yes				
mysqlx_max_connections	Yes	Yes	Yes		Global	Yes
mysqlx_min_worker_threads	Yes	Yes				
Mysqlx_notice_other_sent				Yes	Both	No
Mysqlx_notice_warning_sent				Yes	Both	No
Mysqlx_port				Yes	Global	No
mysqlx_port	Yes	Yes	Yes		Global	No
mysqlx_port_open_timeout	Yes	Yes	Yes		Global	No
Mysqlx_rows_sent				Yes	Both	No
Mysqlx_sessions				Yes	Global	No
Mysqlx_sessions_accepted				Yes	Global	No
Mysqlx_sessions_closed				Yes	Global	No
Mysqlx_sessions_fatal_error				Yes	Global	No
Mysqlx_sessions_killed				Yes	Global	No
Mysqlx_sessions_rejected				Yes	Global	No
Mysqlx_socket				Yes	Global	No
mysqlx-socket	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
mysqlx_socket	Yes	Yes	Yes		Global	No
mysqlx_ssl	Yes	Yes	Yes		Global	No
Mysqlx_ssl_accept_renegotiates				Yes	Global	No
Mysqlx_ssl_accepts				Yes	Global	No
Mysqlx_ssl_active				Yes	Both	No
mysqlx_ssl_ca	Yes	Yes	Yes		Global	No
mysqlx_ssl_capath	Yes	Yes	Yes		Global	No
mysqlx_ssl_cert	Yes	Yes	Yes		Global	No
Mysqlx_ssl_cipher				Yes	Both	No
mysqlx_ssl_cipher	Yes	Yes				
Mysqlx_ssl_cipher_list				Yes	Both	No
mysqlx_ssl_crl	Yes	Yes	Yes		Global	No
mysqlx_ssl_crlpath	Yes	Yes	Yes		Global	No
Mysqlx_ssl_ctx_verify_depth				Yes	Both	No
Mysqlx_ssl_ctx_verify_mode				Yes	Both	No
Mysqlx_ssl_finished_accepts				Yes	Global	No
mysqlx_ssl_key	Yes	Yes	Yes		Global	No
Mysqlx_ssl_server_not_after				Yes	Global	No
Mysqlx_ssl_server_not_before				Yes	Global	No
Mysqlx_ssl_verify_depth				Yes	Global	No
Mysqlx_ssl_verify_mode				Yes	Global	No
Mysqlx_ssl_version				Yes	Both	No
Mysqlx_stmt_create_collection				Yes	Both	No
Mysqlx_stmt_create_collection_index				Yes	Both	No
Mysqlx_stmt_disable_notices				Yes	Both	No
Mysqlx_stmt_drop_collection				Yes	Both	No
Mysqlx_stmt_drop_collection_index				Yes	Both	No
Mysqlx_stmt_enable_notices				Yes	Both	No
Mysqlx_stmt_ensure_collection				Yes	Both	No
Mysqlx_stmt_execute_mysqlx				Yes	Both	No
Mysqlx_stmt_execute_sql				Yes	Both	No
Mysqlx_stmt_execute_xplugin				Yes	Both	No
Mysqlx_stmt_kill_client				Yes	Both	No
Mysqlx_stmt_list_clients				Yes	Both	No
Mysqlx_stmt_list_notices				Yes	Both	No
Mysqlx_stmt_list_objects				Yes	Both	No
Mysqlx_stmt_ping				Yes	Both	No
Mysqlx_worker_threads				Yes	Global	No
Mysqlx_worker_threads_active				Yes	Global	No

20.7.2.2 X Plugin System Variables and Options

The following command options configure X Plugin:

- `--mysqlx_port[=value]`

Command-Line Format		`--mysqlx-port=#`
System Variable	Name	`mysqlx_port`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`33060`
	Min Value	`1`
	Max Value	`65535`

Specifies the port where the X Plugin listens for connections. This option is only read at start up, therefore to change it on a running server requires a restart.

- `--mysqlx_connect_timeout[=value]`

Command-Line Format		`--mysqlx-connect-timeout=#`
System Variable	Name	`mysqlx_connect_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`30`
	Min Value	`1`
	Max Value	`18446744073709551615`

Number of seconds to wait for the first packet to be received from newly connected clients. This option is dynamically changeable.

- `mysqlx_idle_worker_thread_timeout[=value]`

Command-Line Format		`--mysqlx-idle-worker-thread-timeout=#`
Permitted Values	Type	integer
	Default	`60`
	Min Value	`0`
	Max Value	`18446744073709551615`

Time in seconds after which an idle worker thread is terminated.

- `mysqlx_max_allowed_packet[=value]`

Command-Line Format		`--mysqlx-max-allowed-packet=#`
Permitted Values	Type	integer
	Default	`1048576`
	Min Value	`0`
	Max Value	`18446744073709551615`

Maxmum size of a network packet that X Plugin can process. This option is dynamically changeable.

- `mysqlx_max_connections[=value]`

Command-Line Format		`--mysqlx-max-connections=#`
System Variable	Name	`mysqlx_max_connections`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`100`
	Min Value	`1`
	Max Value	`18446744073709551615`

Maximum number of concurrent client connections the X Plugin can accept. This option is dynamically changeable. When modifying this variable, if the new value is smaller than the current number of connections, the new limit is only taken into account for new connections.

- `mysqlx_min_worker_threads[=value]`

Command-Line Format		`--mysqlx-min-worker-threads=#`
Permitted Values	Type	integer
	Default	`2`
	Min Value	`1`
	Max Value	`18446744073709551615`

The minimum number of worker threads the X Plugin uses for handling client requests. The maximum number of worker threads is limited implicitly by `--mysqlx_max_connections`. This option is dynamically changeable.

- `--mysqlx_ssl[=value]`

Command-Line Format		`--mysqlx-ssl=#`
System Variable	Name	`mysqlx_ssl`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	integer

The equivalent of `ssl` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_ca[=value]`

Command-Line Format	`--mysqlx-ssl-ca=file_name`	
System Variable	**Name**	`mysqlx_ssl_ca`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The equivalent of `ssl_ca` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_capath[=value]`

Command-Line Format	`--mysqlx-ssl-capath=dir_name`	
System Variable	**Name**	`mysqlx_ssl_capath`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The equivalent of `ssl_capath` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_cert[=value]`

Command-Line Format	`--mysqlx-ssl-cert`	
System Variable	**Name**	`mysqlx_ssl_key`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The equivalent of `ssl_cert` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_cipher[=value]`

Command-Line Format	`--mysqlx-ssl-cipher=name`	
Permitted Values	**Type**	string

The equivalent of `ssl_cipher` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_crl[=value]`

Command-Line Format	`--mysqlx-ssl-crl=file_name`	
System Variable	**Name**	`mysqlx_ssl_crl`

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The equivalent of `ssl_crl` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_crlpath[=value]`

Command-Line Format	`--mysqlx-ssl-crlpath=directory_name`	
System Variable	**Name**	`mysqlx_ssl_crlpath`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	directory name

The equivalent of `ssl_crlpath` for X Plugin, see that variable for more information.

- `--mysqlx_ssl_key[=value]`

Command-Line Format	`--mysqlx-ssl-key=file_name`	
System Variable	**Name**	`mysqlx_ssl_key`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	file name

The equivalent of `ssl_key` for X Plugin, see that variable for more information.

20.7.3 Monitoring X Plugin

This section describes how to monitor X Plugin. There are two available methods of monitoring, using Performance Schema tables or status variables.

20.7.3.1 Status Variables for X Plugin

The status variables have the following meanings.

- `Mysqlx_bytes_sent`

  The number of bytes sent through the network.

- `Mysqlx_bytes_received`

  The number of bytes received through the network.

- `Mysqlx_stmt_execute_sql`

  The number of StmtExecute requests received for the SQL namespace.

- `Mysqlx_stmt_execute_xplugin`

  The number of StmtExecute requests received for the X Plugin namespace.

- `Mysqlx_stmt_create_collection`

  The number of create collection statements received.

- `Mysqlx_stmt_create_collection_index`

  The number of create collection index statements received.

- `Mysqlx_stmt_drop_collection`

  The number of drop collection statements received.

- `Mysqlx_stmt_drop_collection_index`

  The number of drop collection index statements received.

- `Mysqlx_stmt_list_objects`

  The number of list object statements received.

- `Mysqlx_stmt_enable_notices`

  The number of enable notice statements received.

- `Mysqlx_stmt_disable_notices`

  The number of disable notice statements received.

- `Mysqlx_stmt_list_notices`

  The number of list notice statements received.

- `Mysqlx_stmt_list_clients`

  The number of list client statements received.

- `Mysqlx_stmt_kill_client`

  The number of kill client statements received.

- `Mysqlx_stmt_ping`

  The number of ping statements received.

- `Mysqlx_crud_insert`

  The number of insert requests received.

- `Mysqlx_crud_update`

  The number of update requests received.

- `Mysqlx_crud_delete`

  The number of delete requests received.

- `Mysqlx_crud_find`

  The number of find requests received.

- `Mysqlx_expect_open`

  The number of expectation blocks opened.

- `Mysqlx_expect_close`

The number of expectation blocks closed.

- `Mysqlx_rows_sent`

  The number of rows sent back to clients.

- `Mysqlx_notice_warning_sent`

  The number of warning notices sent back to clients.

- `Mysqlx_notice_other_sent`

  The number of other types of notices sent back to clients.

- `Mysqlx_ssl_cipher`

  The current SSL cipher (empty for non-SSL connections).

- `Mysqlx_ssl_cipher_list`

  A list of possible SSL ciphers (empty for non-SSL connections).

- `Mysqlx_ssl_verify_depth`

  The certificate verification depth for SSL connections.

- `Mysqlx_ssl_verify_mode`

  The certificate verification mode for SSL connections.

- `Mysqlx_ssl_version`

  The name of the protocol used for SSL connections.

- `Mysqlx_sessions`

  The number of sessions that have been opened.

- `Mysqlx_sessions_closed`

  The number of sessions that have been closed.

- `Mysqlx_sessions_fatal_error`

  The number of sessions that have closed with a fatal error.

- `Mysqlx_init_error`

  The number of errors during initialisation.

- `Mysqlx_sessions_accepted`

  The number of session attempts which have been accepted.

- `Mysqlx_sessions_rejected`

  The number of session attempts which have been rejected.

- `Mysqlx_sessions_killed`

  The number of sessions which have been killed.

- `Mysqlx_connections_closed`

The number of connections which have been closed.

- `Mysqlx_connections_accepted`

  The number of connections which have been accepted.

- `Mysqlx_connections_rejected`

  The number of connections which have been rejected.

- `Mysqlx_connection_accept_errors`

  The number of connections which have caused accept errors.

- `Mysqlx_connection_errors`

  The number of connections which have caused errors.

- `Mysqlx_worker_threads`

  The number of worker threads available.

- `Mysqlx_worker_threads_active`

  The number of worker threads currently used.

- `Mysqlx_ssl_active`

  If SSL is active.

- `Mysqlx_ssl_ctx_verify_depth`

  The certificate verification depth limit currently set in ctx.

- `Mysqlx_ssl_ctx_verify_mode`

  The certificate verification mode currently set in ctx.

- `Mysqlx_ssl_finished_accepts`

  The number of successful SSL connections to the server.

- `Mysqlx_ssl_accepts`

  The number of accepted SSL connections.

- `Mysqlx_ssl_server_not_after`

  The last date for which the SSL certificate is valid.

- `Mysqlx_ssl_server_not_before`

  The first date for which the SSL certificate is valid.

- `Mysqlx_ssl_accept_renegotiates`

  The number of negotiations needed to establish the connection.

- `Mysqlx_errors_sent`

  The number of errors sent to clients.

Chapter 21 Group Replication

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This chapter explains MySQL Group Replication and how to install, configure and monitor groups. MySQL Group Replication is a MySQL Server plugin that enables you to create elastic, highly-available, fault-tolerant replication topologies.

Groups can operate in a single-primary mode with automatic primary election, where only one server accepts updates at a time. Alternatively, for more advanced users, groups can be deployed in multi-primary mode, where all servers can accept updates, even if they are issued concurrently.

There is a built-in group membership service that keeps the view of the group consistent and available for all servers at any given point in time. Servers can leave and join the group and the view is updated accordingly. Sometimes servers can leave the group unexpectedly, in which case the failure detection mechanism detects this and notifies the group that the view has changed. This is all automatic.

The chapter is structured as follows:

- Section 21.1, “Group Replication Background” provides an introduction to groups and how Group Replication works.

- Section 21.2, “Getting Started” explains how to configure multiple MySQL Server instances to create a group.
- Section 21.3, “Monitoring Group Replication” explains how to monitor a group.
- Section 21.5, “Group Replication Security” explains how to secure a group.
- Section 21.9, “Group Replication Technical Details” provides in-depth information about how Group Replication works.

21.1 Group Replication Background

This section provides background information on MySQL Group Replication.

The most common way to create a fault-tolerant system is to resort to making components redundant, in other words the component can be removed and the system should continue to operate as expected. This creates a set of challenges that raise complexity of such systems to a whole different level. Specifically, replicated databases have to deal with the fact that they require maintenance and administration of several servers instead of just one. Moreover, as servers are cooperating together to create the group several other classic distributed systems problems have to be dealt with, such as network partitioning or split brain scenarios.

Therefore, the ultimate challenge is to fuse the logic of the database and data replication with the logic of having several servers coordinated in a consistent and simple way. In other words, to have multiple servers agreeing on the state of the system and the data on each and every change that the system goes through. This can be summarized as having servers reaching agreement on each database state transition, so that they all progress as one single database or alternatively that they eventually converge to the same state. Meaning that they need to operate as a (distributed) state machine.

MySQL Group Replication provides distributed state machine replication with strong coordination between servers. Servers coordinate themselves automatically when they are part of the same group. The group can operate in a single-primary mode with automatic primary election, where only one server accepts updates at a time. Alternatively, for more advanced users the group can be deployed in multi-primary mode, where all servers can accept updates, even if they are issued concurrently. This power comes at the expense of applications having to work around the limitations imposed by such deployments.

There is a built-in group membership service that keeps the view of the group consistent and available for all servers at any given point in time. Servers can leave and join the group and the view is updated accordingly. Sometimes servers can leave the group unexpectedly, in which case the failure detection mechanism detects this and notifies the group that the view has changed. This is all automatic.

For a transaction to commit, the majority of the group have to agree on the order of a given transaction in the global sequence of transactions. Deciding to commit or abort a transaction is done by each server individually, but all servers make the same decision. If there is a network partition, resulting in a split where members are unable to reach agreement, then the system does not progress until this issue is resolved. Hence there is also a built-in, automatic, split-brain protection mechanism.

All of this is powered by the provided group communication protocols. These provide a failure detection mechanism, a group membership service, and safe and completely ordered message delivery. All these properties are key to creating a system which ensures that data is consistently replicated across the group of servers. At the very core of this technology lies an implementation of the Paxos algorithm. It acts as the group communication systems engine.

21.1.1 Replication Technologies

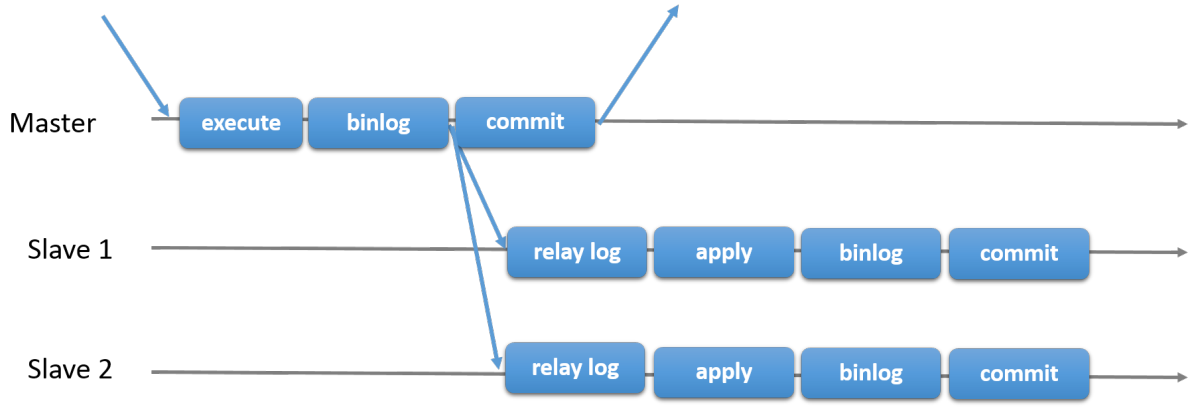
Before getting into the details of MySQL Group Replication, this section introduces some background concepts and an overview of how things work. This provides some context to help understand what

is required for Group Replication and what the differences are between classic asynchronous MySQL Replication and Group Replication.

21.1.1.1 Primary-Secondary Replication

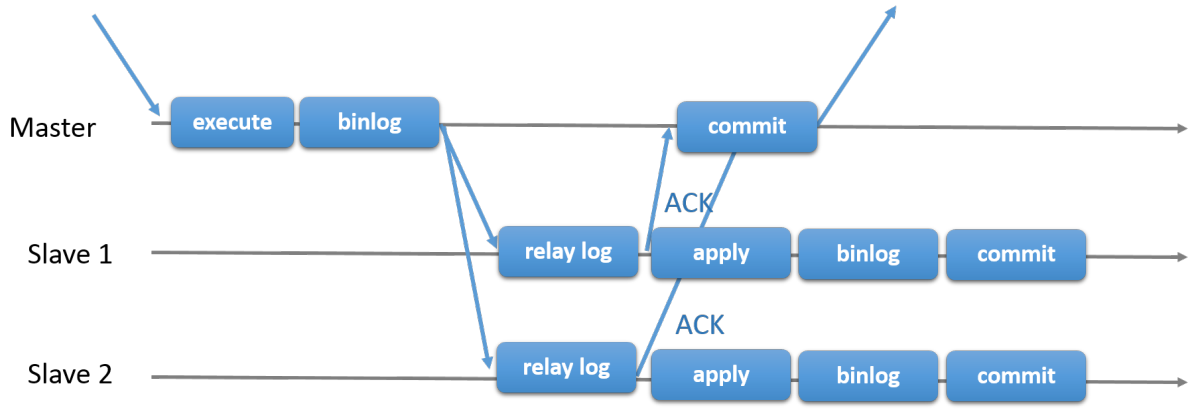
Traditional MySQL Replication provides a simple Primary-Secondary approach to replication. There is a primary (master) and there is one or more secondaries (slaves). The primary executes transactions, commits them and then they are later (thus asynchronously) sent to the secondaries to be either re-executed (in statement-based replication) or applied (in row-based replication). It is a shared-nothing system, where all servers have a full copy of the data by default.

Figure 21.1 MySQL Asynchronous Replication



There is also semisynchronous replication, which adds one synchronization step to the protocol. This means that the Primary waits, at commit time, for the secondary to acknowledge that it has *received* the transaction. Only then does the Primary resume the commit operation.

Figure 21.2 MySQL Semisynchronous Replication



In the two pictures above, you can see a diagram of the classic asynchronous MySQL Replication protocol (and its semisynchronous variant as well). Blue arrows represent messages exchanged between servers or messages exchanged between servers and the client application.

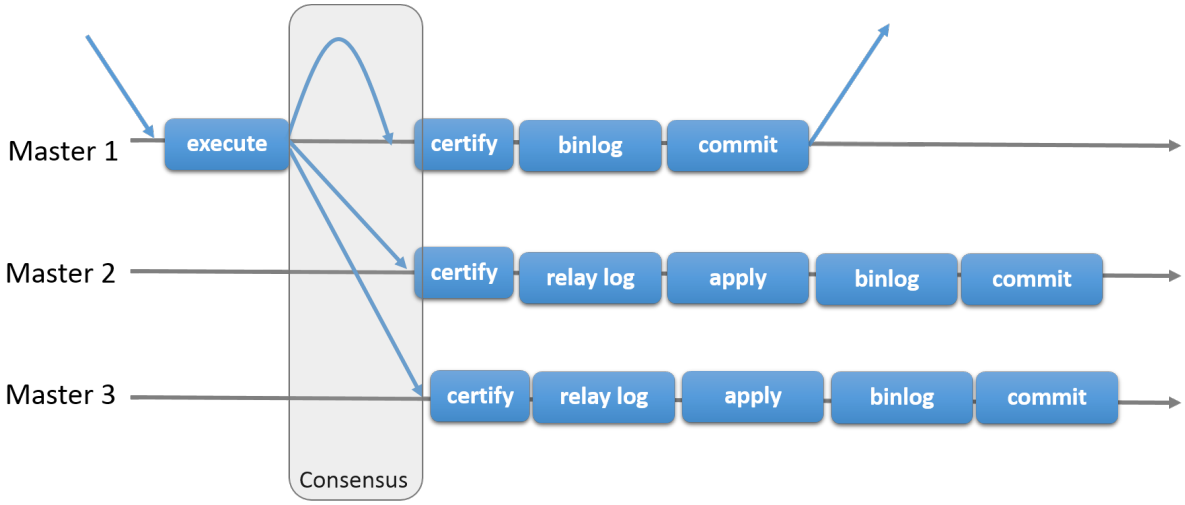
21.1.1.2 Group Replication

Group Replication is a technique that can be used to implement fault-tolerant systems. The replication group is a set of servers that interact with each other through message passing. The communication layer provides a set of guarantees such as atomic message and total order message delivery. These are very powerful properties that translate into very useful abstractions that one can resort to build more advanced database replication solutions.

MySQL Group Replication builds on top of such properties and abstractions and implements a multi-master update everywhere replication protocol. In essence, a replication group is formed by multiple servers and each server in the group may execute transactions independently. But all read-write (RW) transactions commit only after they have been approved by the group. Read-only (RO) transactions need no coordination within the group and thus commit immediately. In other words, for any RW transaction the group needs to decide whether it commits or not, thus the commit operation is not a unilateral decision from the originating server. To be precise, when a transaction is ready to commit at the originating server, the server atomically broadcasts the write values (rows changed) and the correspondent write set (unique identifiers of the rows that were updated). Then a global total order is established for that transaction. Ultimately, this means that all servers receive the same set of transactions in the same order. As a consequence, all servers apply the same set of changes in the same order, therefore they remain consistent within the group.

However, there may be conflicts between transactions that execute concurrently on different servers. Such conflicts are detected by inspecting the write sets of two different and concurrent transactions, in a process called *certification*. If two concurrent transactions, that executed on different servers, update the same row, then there is a conflict. The resolution procedure states that the transaction that was ordered first commits on all servers, whereas the transaction ordered second aborts, and thus is rolled back on the originating server and dropped by the other servers in the group. This is in fact a distributed first commit wins rule.

Figure 21.3 MySQL Group Replication Protocol



Finally, Group Replication is a shared-nothing replication scheme where each server has its own entire copy of the data.

The Figure above depicts the MySQL Group Replication protocol and by comparing it to MySQL Replication (or even MySQL semisynchronous replication) you can see some differences. Note that some underlying consensus and Paxos related messages are missing from this picture for the sake of clarity.

21.1.2 Group Replication Use Cases

Group Replication enables you to create fault-tolerant systems with redundancy by replicating the system state throughout a set of servers. Consequently, even if some of the servers fail, as long it is not all or a majority, the system is still available, and all it could have degraded performance or scalability, it is still available. Server failures are isolated and independent. They are tracked by a group membership service which relies on a distributed failure detector that is able to signal when any servers leave the group, either voluntarily or due to an unexpected halt. There is a distributed recovery procedure to ensure that when servers join the group they are brought up to date automatically. There is no need for server fail-over, and the multi-master update everywhere nature ensures that not

even updates are blocked in the event of a single server failure. Therefore MySQL Group Replication guarantees that the database service is continuously available.

It is important to understand that although the database service is available, in the event of a server crash, those clients connected to it must be redirected, or failed over, to a different server. This is not something Group Replication attempts to resolve. A connector, load balancer, router, or some form of middleware are more suitable to deal with this issue.

To summarize, MySQL Group Replication provides a highly available, highly elastic, dependable MySQL service.

21.1.2.1 Examples of Use Case Scenarios

The following examples are typical use cases for Group Replication.

- *Elastic Replication* - Environments that require a very fluid replication infrastructure, where the number of servers has to grow or shrink dynamically and with as few side-effects as possible. For instance, database services for the cloud.
- *Highly Available Shards* - Sharding is a popular approach to achieve write scale-out. Use MySQL Group Replication to implement highly available shards, where each shard maps to a replication group.
- *Alternative to Master-Slave replication* - In certain situations, using a single master server makes it a single point of contention. Writing to an entire group may prove more scalable under certain circumstances.
- *Autonomic Systems* - Additionally, you can deploy MySQL Group Replication purely for the automation that is built into the replication protocol (described already in this and previous chapters).

21.1.3 Group Replication Details

This section presents details about some of the services that Group Replication builds on.

21.1.3.1 Failure Detection

There is a failure detection mechanism provided that is able to find and report which servers are silent and as such assumed to be dead. At a high level, the failure detector is a distributed service that provides information about which servers may be dead (suspicions). Later if the group agrees that the suspicions are probably true, then the group decides that a given server has indeed failed. This means that the remaining members in the group take a coordinated decision to exclude a given member.

Suspicions are triggered when servers go mute. When server A does not receive messages from server B during a given period, a timeout occurs and a suspicion is raised.

If a server gets isolated from the rest of the group, then it suspects that all others have failed. Being unable to secure agreement with the group (as it cannot secure a quorum), its suspicion does not have consequences. When a server is isolated from the group in this way, it is unable to execute any local transactions.

21.1.3.2 Group Membership

MySQL Group Replication relies on a group membership service. This is built into the plugin. It defines which servers are online and participating in the group. The list of online servers is often referred to as a *view*. Therefore, every server in the group has a consistent view of which are the members participating actively in the group at a given moment in time.

Servers have to agree not only on transaction commits, but also which is the current view. Therefore, if servers agree that a new server becomes part of the group, then the group itself is reconfigured to

integrate that server in it, triggering a view change. The opposite also happens, if a server leaves the group, voluntarily or not, then the group dynamically rearranges its configuration and a view change is triggered.

Note though that when a member leaves voluntarily, it first initiates a dynamic group reconfiguration. This triggers a procedure, where all members have to agree on the new view without the leaving server. However, if a member leaves involuntarily (for example it has stopped unexpectedly or the network connection is down) then the failure detection mechanism realizes this fact and a reconfiguration of the group is proposed, this one without the failed member. As mentioned this requires agreement from the majority of servers in the group. If the group is not able to reach agreement (for example it partitioned in such a way that there is no majority of servers online), then the system is not be able to dynamically change the configuration and as such, blocks to prevent a split-brain situation. Ultimately, this means that the administrator needs to step in and fix this.

21.1.3.3 Fault-tolerance

MySQL Group Replication builds on an implementation of the Paxos distributed algorithm to provide distributed coordination between servers. As such, it requires a majority of servers to be active to reach quorum and thus make a decision. This has direct impact on the number of failures the system can tolerate without compromising itself and its overall functionality. The number of servers (n) needed to tolerate `f` failures is then `n = 2 x f + 1`.

In practice this means that to tolerate one failure the group must have three servers in it. As such if one server fails, there are still two servers to form a majority (two out of three) and allow the system to continue to make decisions automatically and progress. However, if a second server fails *involuntarily*, then the group (with one server left) blocks, because there is no majority to reach a decision.

The following is a small table illustrating the formula above.

Group Size	Majority	Instant Failures Tolerated
1	1	0
2	2	0
3	2	1
4	3	1
5	3	2
6	4	2
7	4	3

The next Chapter covers technical aspects of Group Replication.

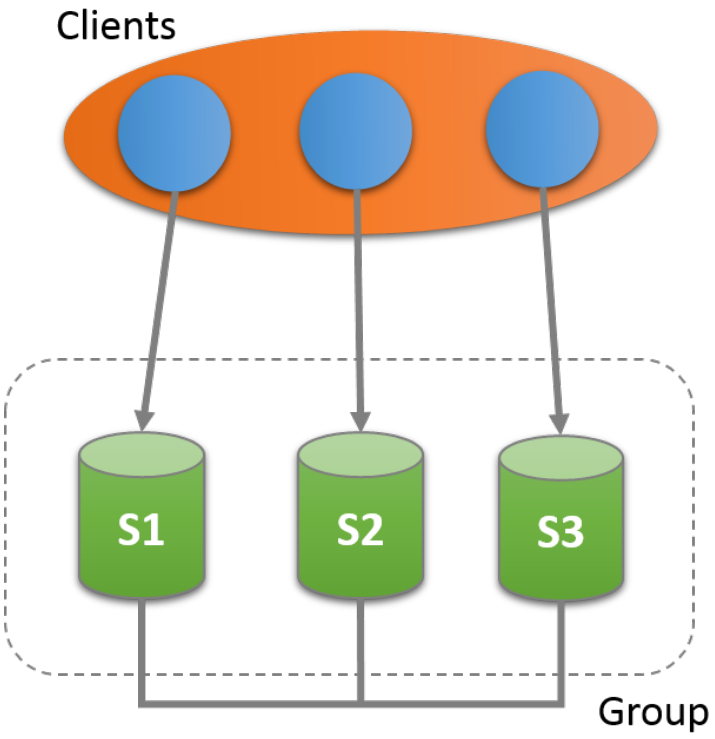
21.2 Getting Started

MySQL Group Replication is provided as a plugin to MySQL server, and each server in a group requires configuration and installation of the plugin. This section provides a detailed tutorial with the steps required to create a replication group with at least three servers.

21.2.1 Deploying Group Replication in Single-Primary Mode

Each of the server instances in a group can run on an independent physical machine, or on the same machine. This section explains how to create a replication group with three MySQL Server instances on one physical machine. This means that three data directories are needed, one per server instance, and that you need to configure each instance independently.

Figure 21.4 Group Architecture



This tutorial explains how to get and deploy MySQL Server with the Group Replication plugin, how to configure each server instance before creating a group, and how to use Performance Schema monitoring to verify that everything is working correctly.

21.2.1.1 Deploying Instances for Group Replication

The first step is to deploy three instances of MySQL Server. Group Replication is a built-in MySQL plugin provided with MySQL Server 5.7.17 and later. For more background information on MySQL plugins, see Section 5.6, "MySQL Server Plugins". This procedure assumes that MySQL Server was downloaded and unpacked into the directory named `mysql-5.7`. The following procedure uses one physical machine, therefore each MySQL server instance requires a specific data directory for the instance. Create the data directories in a directory named `data` and initialize each one.

```
mkdir data
mysql-5.7/bin/mysqld --initialize-insecure --basedir=$PWD/mysql-5.7 --datadir=$PWD/data/s1
mysql-5.7/bin/mysqld --initialize-insecure --basedir=$PWD/mysql-5.7 --datadir=$PWD/data/s2
mysql-5.7/bin/mysqld --initialize-insecure --basedir=$PWD/mysql-5.7 --datadir=$PWD/data/s3
```

Inside `data/s1`, `data/s2`, `data/s3` is an initialized data directory, containing the mysql system database and related tables and much more. To learn more about the initialization procedure, see Section 2.9.1.1, "Initializing the Data Directory Manually Using mysqld".

Warning

Do not use `--initialize-insecure` in production environments, it is only used here to simplify the tutorial. For more information on security settings, see Section 21.5, "Group Replication Security".

21.2.1.2 Configuring an Instance for Group Replication

This section explains the configuration settings required for MySQL Server instances that you want to use for Group Replication. For background information, see Section 21.7.2, "Group Replication Limitations".

Group Replication Server Settings

To install and use the Group Replication plugin you must configure the MySQL Server instance correctly. It is recommended to store the configuration in your `my.cnf` file. See Section 4.2.6, "Using Option Files" for more information. Unless stated otherwise, what follows is the configuration for the first instance in the group, referred to as *s1* in this procedure. The following section shows an example server configuration.

```
[mysqld]

# server configuration
datadir=<full_path_to_data>/data/s1
basedir=<full_path_to_bin>/mysql-5.7/

port=24801
socket=<full_path_to_sock_dir>/s1.sock
```

These settings configure MySQL server to use the data directory created earlier and which port the server should open and start listening for incoming connections.

Note

The non-default port of 24801 is used because in this tutorial the three server instances use the same hostname. In a setup with three different machines this would not be required.

Replication Framework

The following settings configure replication according to the MySQL Group Replication requirements.

```
server_id=1
gtid_mode=ON
enforce_gtid_consistency=ON
master_info_repository=TABLE
relay_log_info_repository=TABLE
binlog_checksum=NONE
log_slave_updates=ON
log_bin=binlog
binlog_format=ROW
```

These settings configure the server to use the unique identifier number 1, to enable global transaction identifiers and to store replication metadata in system tables instead of files. Additionally, it instructs the server to turn on binary logging, use row-based format and disable binary log event checksums. For more details see Section 21.7.2, "Group Replication Limitations".

Group Replication Settings

At this point the `my.cnf` file ensures that the server is configured and is instructed to instantiate the replication infrastructure under a given configuration. The following section configures the Group Replication settings for the server.

```
transaction_write_set_extraction=XXHASH64
loose-group_replication_group_name="aaaaaaaa-aaaa-aaaa-aaaa-aaaaaaaaaaaa"
loose-group_replication_start_on_boot=off
loose-group_replication_local_address= "127.0.0.1:24901"
loose-group_replication_group_seeds= "127.0.0.1:24901,127.0.0.1:24902,127.0.0.1:24903"
loose-group_replication_bootstrap_group= off
```

Note

The `loose-` prefix used for the group_replication variables above instructs the server to continue to start if the Group Replication plugin has not been loaded at the time the server is started.

- *Line 1* instructs the server that for each transaction it has to collect the write set and encode it as a hash using the XXHASH64 hashing algorithm.
- *Line 2* tells the plugin that the group that it is joining, or creating, is named "aaaaaaaa-aaaa-aaaa-aaaa-aaaaaaaaaaaa".
- *Line 3* instructs the plugin to not start operations automatically when the server starts.
- *Line 4* tells the plugin to use the IP address 127.0.0.1, or localhost, and port 24901 for incoming connections from other members in the group.

Important

The server listens on this port for member-to-member connections. This port must not be used for user applications at all, it must be reserved for communication between different members of the group while running Group Replication.

The local address configured by `group_replication_local_address` must be accessible to all group members. For example, if each server instance is on a different machine use the IP and port of the machine, such as 10.0.0.1:33061. The recommended port for `group_replication_local_address` is 33061, but in this tutorial we use three server instances running on one machine, thus ports 24901 to 24903 are used.

- *Line 5* tells the plugin that the following members on those hosts and ports should be contacted in case it needs to join the group. These are seed members, which are used when this member wants to connect to the group. Upon joining the server contacts one of them first (seed) and then it asks the group to reconfigure to allow the joiner to be accepted in the group. Note that this option does *not* need to list all members in the group, but rather a list of servers that should be contacted in case this server wishes to join the group.

  The server that starts the group does not make use of this option, since it is the initial server and as such, it is in charge of bootstrapping the group. The second server joining asks the one and only member in the group to join and then the group expands. The third server joining can ask any of these two to join, and then the group expands again. Subsequent servers repeat this procedure when joining.

Warning

When joining multiple servers at the same time, make sure that they point to seed members that are already in the group. Do not use members that are also joining the group as as seeds, because they may not yet be in the group when contacted.

It is good practice to start the bootstrap member first, and let it create the group. Then make it the seed member for the rest of the members that are joining. This ensures that there is a group formed when joining the rest of the members.

Creating a group and joining multiple members at the same time is not supported. It may work, but chances are that the operations race and then the act of joining the group ends up in an error or a time out.

- *Line 6* instructs the plugin whether to boostrap the group or not.

Important

This option must only be used on one server instance at any time, usually the first time you bootstrap the group (or in case the entire group is brought down and back up again). If you bootstrap the group multiple times, for example when multiple server instances have this option set, then they could create an

artificial split brain scenario, in which two distinct groups with the same name exist. Disable this option after the first server instance comes online.

Configuration for all servers in the group is quite similar. You need to change the specifics about each server (for example `server_id`, `datadir`, `group_replication_local_address`). This is illustrated later in this tutorial.

21.2.1.3 User Credentials

Group Replication uses the asynchronous replication protocol to achieve distributed recovery, synchronizing group members before joining them to the group. The distributed recovery process relies on a replication channel named `group_replication_recovery` which is used to transfer transactions between group members. Therefore you need to set up a replication user with the correct permissions so that Group Replication can establish direct member-to-member recovery replication channels.

Start the server:

```
mysql-5.7/bin/mysqld --defaults-file=data/s1/s1.cnf
```

Create a MySQL user with the `REPLICATION-SLAVE` privilege. This process should *not* be captured in the binary log to avoid the changes being propagated to other server instances. In the following example the user *`rpl_user`* with the password *`rpl_pass`* is shown. When configuring your servers use a suitable user name and password. Connect to server s1 and issue the following statements:

```
mysql> SET SQL_LOG_BIN=0;
Query OK, 0 rows affected (0,00 sec)

mysql> CREATE USER rpl_user@'%' IDENTIFIED BY 'rpl_pass';
Query OK, 0 rows affected (0,00 sec)

mysql> GRANT REPLICATION SLAVE ON *.* TO rpl_user@'%';
Query OK, 0 rows affected, 1 warning (0,00 sec)

mysql> FLUSH PRIVILEGES;
Query OK, 0 rows affected (0,00 sec)

mysql> SET SQL_LOG_BIN=1;
Query OK, 0 rows affected (0,00 sec)
```

Once the user has been configured as above, use the `CHANGE MASTER TO` statement to configure the server to use the given credentials for the `group_replication_recovery` replication channel the next time it needs to recover its state from another member. Issue the following, replacing *`rpl_user`* and *`rpl_pass`* with the values used when creating the user.

```
mysql> CHANGE MASTER TO MASTER_USER='rpl_user', MASTER_PASSWORD='rpl_pass' \\
       FOR CHANNEL 'group_replication_recovery';
Query OK, 0 rows affected, 2 warnings (0,01 sec)
```

Distributed recovery is the first step taken by a server that joins the group. If these credentials are not set correctly as shown, the server cannot run the recovery process and gain synchrony with the other group members, and hence ultimately cannot join the group. Similarly, if the member cannot correctly identify the other members via the server's `hostname` the recovery process can fail. It is recommended that operating systems running MySQL have a properly configured unique `hostname`, either using DNS or local settings. This `hostname` can be verified in the `Member_host` column of the `performance_schema.replication_group_members` table. If multiple group members externalize a default `hostname` set by the operating system, there is a chance of the member not resolving to the correct member address and not being able to join the group. In such a situation use `report_host` to configure a unique `hostname` to be externalized by each of the servers.

21.2.1.4 Launching Group Replication

Once server s1 has been configured and started, install the Group Replication plugin. Connect to the server and issue the following command:

```
INSTALL PLUGIN group_replication SONAME 'group_replication.so';
```

Important

The `mysql.session` user must exist before you can load Group Replication. `mysql.session` was added in MySQL version 8.0.2. If your data dictionary was initialized using an earlier version you must run the `mysql_upgrade` procedure. If the upgrade is not run, Group Replication fails to start with the error message `There was an error when trying to access the server with user: mysql.session_user@localhost. Make sure the user is present in the server and that mysql_upgrade was ran after a server update.`.

To check that the plugin was installed successfully, issue `SHOW PLUGINS;` and check the output. It should show something like this:

```
mysql> SHOW PLUGINS;
+----------------------------+----------+--------------------+----------------------+-------------+
| Name                       | Status   | Type               | Library              | License     |
+----------------------------+----------+--------------------+----------------------+-------------+
| binlog                     | ACTIVE   | STORAGE ENGINE     | NULL                 | PROPRIETARY |

(...)

| group_replication          | ACTIVE   | GROUP REPLICATION  | group_replication.so | PROPRIETARY |
+----------------------------+----------+--------------------+----------------------+-------------+
```

To start the group, instruct server s1 to bootstrap the group and then start Group Replication. This bootstrap should only be done by a single server, the one that starts the group and only once. This is why the value of the bootstrap configuration option was not saved in the configuration file. If it is saved in the configuration file, upon restart the server automatically bootstraps a second group with the same name. This would result in two distinct groups with the same name. The same reasoning applies to stopping and restarting the plugin with this option set to `ON`.

```
SET GLOBAL group_replication_bootstrap_group=ON;
START GROUP_REPLICATION;
SET GLOBAL group_replication_bootstrap_group=OFF;
```

Once the `START GROUP_REPLICATION` statement returns, the group has been started. You can check that the group is now created and that there is one member in it:

```
mysql> SELECT * FROM performance_schema.replication_group_members;
+----------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME               | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+----------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | ce9be252-2b71-11e6-b8f4-00212844f856 | myhost      |       24801 | ONLINE
+----------------------------+--------------------------------------+-------------+-------------+-------
1 row in set (0,00 sec)
```

The information in this table confirms that there is a member in the group with the unique identifier `ce9be252-2b71-11e6-b8f4-00212844f856`, that it is `ONLINE` and is at `myhost` listening for client connections on port `24801`.

For the purpose of demonstrating that the server is indeed in a group and that it is able to handle load, create a table and add some content to it.

```
mysql> CREATE DATABASE test;
Query OK, 1 row affected (0,00 sec)

mysql> use test
Database changed
mysql> CREATE TABLE t1 (c1 INT PRIMARY KEY, c2 TEXT NOT NULL);
Query OK, 0 rows affected (0,00 sec)

mysql> INSERT INTO t1 VALUES (1, 'Luis');
Query OK, 1 row affected (0,01 sec)
```

Check the content of table `t1` and the binary log.

```
mysql> SELECT * FROM t1;
+----+------+
| c1 | c2   |
+----+------+
|  1 | Luis |
+----+------+
1 row in set (0,00 sec)

mysql> SHOW BINLOG EVENTS;
+---------------+-----+----------------+-----------+-------------+-----------------------------------------
| Log_name      | Pos | Event_type     | Server_id | End_log_pos | Info
+---------------+-----+----------------+-----------+-------------+-----------------------------------------
| binlog.000001 |   4 | Format_desc    |         1 |         123 | Server ver: 5.7.17-gr080-log, Binlog ver
| binlog.000001 | 123 | Previous_gtids |         1 |         150 |
| binlog.000001 | 150 | Gtid           |         1 |         211 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa-
| binlog.000001 | 211 | Query          |         1 |         270 | BEGIN
| binlog.000001 | 270 | View_change    |         1 |         369 | view_id=14724817264259180:1
| binlog.000001 | 369 | Query          |         1 |         434 | COMMIT
| binlog.000001 | 434 | Gtid           |         1 |         495 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa-
| binlog.000001 | 495 | Query          |         1 |         585 | CREATE DATABASE test
| binlog.000001 | 585 | Gtid           |         1 |         646 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa-
| binlog.000001 | 646 | Query          |         1 |         770 | use `test`; CREATE TABLE t1 (c1 INT PRIM
| binlog.000001 | 770 | Gtid           |         1 |         831 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa-
| binlog.000001 | 831 | Query          |         1 |         899 | BEGIN
| binlog.000001 | 899 | Table_map      |         1 |         942 | table_id: 108 (test.t1)
| binlog.000001 | 942 | Write_rows     |         1 |         984 | table_id: 108 flags: STMT_END_F
| binlog.000001 | 984 | Xid            |         1 |        1011 | COMMIT /* xid=38 */
+---------------+-----+----------------+-----------+-------------+-----------------------------------------
15 rows in set (0,00 sec)
```

As seen above, the database and the table objects were created and their corresponding DDL statements were written to the binary log. Also, the data was inserted into the table and written to the binary log. The importance of the binary log entries is illustrated in the following section when the group grows and distributed recovery is executed as new members try to catch up and become online.

21.2.1.5 Adding Instances to the Group

At this point, the group has one member in it, server s1, which has some data in it. It is now time to expand the group by adding the other two servers configured previously.

Adding a Second Instance

In order to add a second instance, server s2, first create the configuration file for it. The configuration is similar to the one used for server s1, except for things such as the location of the data directory, the ports that s2 is going to be listening on or its `server_id`. These different lines are highlighted in the listing below.

```
[mysqld]

# server configuration
datadir=<full_path_to_data>/data/s2
basedir=<full_path_to_bin>/mysql-5.7/
```

```
port=24802
socket=<full_path_to_sock_dir>/s2.sock

#
# Replication configuration parameters
#
server_id=2
gtid_mode=ON
enforce_gtid_consistency=ON
master_info_repository=TABLE
relay_log_info_repository=TABLE
binlog_checksum=NONE
log_slave_updates=ON
log_bin=binlog
binlog_format=ROW

#
# Group Replication configuration
#
transaction_write_set_extraction=XXHASH64
loose-group_replication_group_name="aaaaaaaa-aaaa-aaaa-aaaa-aaaaaaaaaaaa"
loose-group_replication_start_on_boot=off
loose-group_replication_local_address= "127.0.0.1:24902"
loose-group_replication_group_seeds= "127.0.0.1:24901,127.0.0.1:24902,127.0.0.1:24903"
loose-group_replication_bootstrap_group= off
```

Similar to the procedure for server s1, with the configuration file in place you launch the server.

```
mysql-5.7/bin/mysqld --defaults-file=data/s2/s2.cnf
```

Then configure the recovery credentials as follows. The commands are the same as used when setting up server s1 as the user is shared within the group. Issue the following statements on s2.

```
SET SQL_LOG_BIN=0;
CREATE USER rpl_user@'%';
GRANT REPLICATION SLAVE ON *.* TO rpl_user@'%' IDENTIFIED BY 'rpl_pass';
SET SQL_LOG_BIN=1;
CHANGE MASTER TO MASTER_USER='rpl_user', MASTER_PASSWORD='rpl_pass' \\
 FOR CHANNEL 'group_replication_recovery';
```

Install the Group Replication plugin and start the process of joining the server to the group. The following example installs the plugin in the same way as used while deploying server s1.

```
mysql> INSTALL PLUGIN group_replication SONAME 'group_replication.so';
Query OK, 0 rows affected (0,01 sec)
```

Add server s2 to the group.

```
mysql> START GROUP_REPLICATION;
Query OK, 0 rows affected (44,88 sec)
```

Unlike the previous steps that were the same as those executed on s1, here there is a difference in that you do *not* issue `SET GLOBAL group_replication_bootstrap_group=ON;` before starting Group Replication, because the group has already been created and bootstrapped by server s1. At this point server s2 only needs to be added to the already existing group.

Important

When Group Replication starts successfully and the server joins the group, if the `super_read_only` variable is set to 1 then the `super_read_only` variable is set to 0. By setting `super_read_only` to 1 in the member's `.cfg` file, you can ensure that servers which fail when starting Group Replication for any reason do not accept transactions.

Checking the `performance_schema.replication_group_members` table again shows that there are now two *ONLINE* servers in the group.

```
mysql> SELECT * FROM performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-----------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER_STA
+---------------------------+--------------------------------------+-------------+-------------+-----------
| group_replication_applier | 395409e1-6dfa-11e6-970b-00212844f856 | myhost      |       24801 | ONLINE
| group_replication_applier | ac39f1e6-6dfa-11e6-a69d-00212844f856 | myhost      |       24802 | ONLINE
+---------------------------+--------------------------------------+-------------+-------------+-----------
2 rows in set (0,00 sec)
```

As server s2 is also marked as ONLINE, it must have already caught up with server s1 automatically. Verify that it has indeed synchronized with server s1 as follows.

```
mysql> SHOW DATABASES LIKE 'test';
+-----------------+
| Database (test) |
+-----------------+
| test            |
+-----------------+
1 row in set (0,00 sec)

mysql> SELECT * FROM test.t1;
+----+------+
| c1 | c2   |
+----+------+
|  1 | Luis |
+----+------+
1 row in set (0,00 sec)

mysql> SHOW BINLOG EVENTS;
+---------------+------+----------------+-----------+-------------+-------------------------------------
| Log_name      | Pos  | Event_type     | Server_id | End_log_pos | Info
+---------------+------+----------------+-----------+-------------+-------------------------------------
| binlog.000001 |    4 | Format_desc    |         2 |         123 | Server ver: 5.7.17-log, Binlog ver: 4
| binlog.000001 |  123 | Previous_gtids |         2 |         150 |
| binlog.000001 |  150 | Gtid           |         1 |         211 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  211 | Query          |         1 |         270 | BEGIN
| binlog.000001 |  270 | View_change    |         1 |         369 | view_id=14724832985483517:1
| binlog.000001 |  369 | Query          |         1 |         434 | COMMIT
| binlog.000001 |  434 | Gtid           |         1 |         495 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  495 | Query          |         1 |         585 | CREATE DATABASE test
| binlog.000001 |  585 | Gtid           |         1 |         646 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  646 | Query          |         1 |         770 | use `test`; CREATE TABLE t1 (c1 INT PRI
| binlog.000001 |  770 | Gtid           |         1 |         831 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  831 | Query          |         1 |         890 | BEGIN
| binlog.000001 |  890 | Table_map      |         1 |         933 | table_id: 108 (test.t1)
| binlog.000001 |  933 | Write_rows     |         1 |         975 | table_id: 108 flags: STMT_END_F
| binlog.000001 |  975 | Xid            |         1 |        1002 | COMMIT /* xid=30 */
| binlog.000001 | 1002 | Gtid           |         1 |        1063 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 | 1063 | Query          |         1 |        1122 | BEGIN
| binlog.000001 | 1122 | View_change    |         1 |        1261 | view_id=14724832985483517:2
| binlog.000001 | 1261 | Query          |         1 |        1326 | COMMIT
+---------------+------+----------------+-----------+-------------+-------------------------------------
19 rows in set (0,00 sec)
```

As seen above, the second server has been added to the group and it has replicated the changes from server s1 automatically. According to the distributed recovery procedure, this means that just after joining the group and immediately before being declared online, server s2 has connected to server s1 automatically and fetched the missing data from it. In other words, it copied transactions from the binary log of s1 that it was missing, up to the point in time that it joined the group.

Adding Additional Instances

Adding additional instances to the group is essentially the same sequence of steps as adding the second server, except that the configuration has to be changed as it had to be for server s2. To summarise the required commands:

1. Create the configuration file

```
[mysqld]

# server configuration
datadir=<full_path_to_data>/data/s3
basedir=<full_path_to_bin>/mysql-5.7/

port=24803
socket=<full_path_to_sock_dir>/s3.sock

#
# Replication configuration parameters
#
server_id=3
gtid_mode=ON
enforce_gtid_consistency=ON
master_info_repository=TABLE
relay_log_info_repository=TABLE
binlog_checksum=NONE
log_slave_updates=ON
log_bin=binlog
binlog_format=ROW

#
# Group Replication configuration
#
transaction_write_set_extraction=XXHASH64
loose-group_replication_group_name="aaaaaaaa-aaaa-aaaa-aaaa-aaaaaaaaaaaa"
loose-group_replication_start_on_boot=off
loose-group_replication_local_address= "127.0.0.1:24903"
loose-group_replication_group_seeds= "127.0.0.1:24901,127.0.0.1:24902,127.0.0.1:24903"
loose-group_replication_bootstrap_group= off
```

2. Start the server

```
mysql-5.7/bin/mysqld --defaults-file=data/s3/s3.cnf
```

3. Configure the recovery credentials for the group_replication_recovery channel.

```
SET SQL_LOG_BIN=0;
CREATE USER rpl_user@'%';
GRANT REPLICATION SLAVE ON *.* TO rpl_user@'%' IDENTIFIED BY 'rpl_pass';
FLUSH PRIVILEGES;
SET SQL_LOG_BIN=1;
CHANGE MASTER TO MASTER_USER='rpl_user', MASTER_PASSWORD='rpl_pass'  \\
FOR CHANNEL 'group_replication_recovery';
```

4. Install the Group Replication plugin and start it.

```
INSTALL PLUGIN group_replication SONAME 'group_replication.so';
START GROUP_REPLICATION;
```

At this point server s3 is booted and running, has joined the group and caught up with the other servers in the group. Consulting the `performance_schema.replication_group_members` table again confirms this is the case.

```
mysql> SELECT * FROM performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+---------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | 395409e1-6dfa-11e6-970b-00212844f856 | myhost      |       24801 | ONLINE
| group_replication_applier | 7eb217ff-6df3-11e6-966c-00212844f856 | myhost      |       24803 | ONLINE
| group_replication_applier | ac39f1e6-6dfa-11e6-a69d-00212844f856 | myhost      |       24802 | ONLINE
+---------------------------+--------------------------------------+-------------+-------------+-------
3 rows in set (0,00 sec)
```

Issuing this same query on server s2 or server s1 yields the same result. Also, you can verify that server s3 has also caught up:

```
mysql> SHOW DATABASES LIKE 'test';
+-----------------+
| Database (test) |
+-----------------+
| test            |
+-----------------+
1 row in set (0,00 sec)

mysql> SELECT * FROM test.t1;
+----+------+
| c1 | c2   |
+----+------+
|  1 | Luis |
+----+------+
1 row in set (0,00 sec)

mysql>  SHOW BINLOG EVENTS;
+---------------+------+----------------+-----------+-------------+-------------------------------------
| Log_name      | Pos  | Event_type     | Server_id | End_log_pos | Info
+---------------+------+----------------+-----------+-------------+-------------------------------------
| binlog.000001 |    4 | Format_desc    |         3 |         123 | Server ver: 5.7.17-log, Binlog ver: 4
| binlog.000001 |  123 | Previous_gtids |         3 |         150 |
| binlog.000001 |  150 | Gtid           |         1 |         211 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  211 | Query          |         1 |         270 | BEGIN
| binlog.000001 |  270 | View_change    |         1 |         369 | view_id=14724832985483517:1
| binlog.000001 |  369 | Query          |         1 |         434 | COMMIT
| binlog.000001 |  434 | Gtid           |         1 |         495 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  495 | Query          |         1 |         585 | CREATE DATABASE test
| binlog.000001 |  585 | Gtid           |         1 |         646 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  646 | Query          |         1 |         770 | use `test`; CREATE TABLE t1 (c1 INT PRI
| binlog.000001 |  770 | Gtid           |         1 |         831 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 |  831 | Query          |         1 |         890 | BEGIN
| binlog.000001 |  890 | Table_map      |         1 |         933 | table_id: 108 (test.t1)
| binlog.000001 |  933 | Write_rows     |         1 |         975 | table_id: 108 flags: STMT_END_F
| binlog.000001 |  975 | Xid            |         1 |        1002 | COMMIT /* xid=29 */
| binlog.000001 | 1002 | Gtid           |         1 |        1063 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 | 1063 | Query          |         1 |        1122 | BEGIN
| binlog.000001 | 1122 | View_change    |         1 |        1261 | view_id=14724832985483517:2
| binlog.000001 | 1261 | Query          |         1 |        1326 | COMMIT
| binlog.000001 | 1326 | Gtid           |         1 |        1387 | SET @@SESSION.GTID_NEXT= 'aaaaaaaa-aaaa
| binlog.000001 | 1387 | Query          |         1 |        1446 | BEGIN
| binlog.000001 | 1446 | View_change    |         1 |        1585 | view_id=14724832985483517:3
| binlog.000001 | 1585 | Query          |         1 |        1650 | COMMIT
+---------------+------+----------------+-----------+-------------+-------------------------------------
23 rows in set (0,00 sec)
```

21.3 Monitoring Group Replication

Use the Perfomance Schema tables to monitor Group Replication, assuming that the Performance Schema is enabled. Group Replication adds the following tables:

- `performance_schema.replication_group_member_stats`
- `performance_schema.replication_group_members`

These existing Perfomance Schema replication tables also show information regarding Group Replication.

- `performance_schema.replication_connection_status`
- `performance_schema.replication_applier_status`

The replication channels created by the Group Replication plugin are named:

- `group_replication_recovery` - This channel is used for the replication changes that are related to the distributed recovery phase.

- `group_replication_applier` - This channel is used for the incoming changes from the group. This is the channel used to apply transactions coming directly from the group.

The following sections describe the information available in each of these tables.

21.3.1 Replication_group_member_stats

Each member in a replication group certifies and applies transactions committed by the group. Statistics regarding the certifier and applier procedures are useful to understand how the applier queue is growing, how many conflicts have been found, how many transactions were checked, which transactions are committed everywhere, and so on. The table `performance_schema.replication_group_member_stats` provides the following information related to the certification process.

Table 21.1 replication_group_member_stats

Field	Description
Channel_name	The name of the Group Replication channel.
Member_id	This gives the member server UUID to which we are currently connected to. This has a different value for each member in the group. This also serves a key since it is unique to each member.
Count_Transactions_in_queue	Number of transactions in queue pending conflict detection checks. Once checked for conflicts, and if they pass the check, they are queued to be applied as well.
Count_transactions_checked	Indicates the number of transactions that were checked for conflicts.
Count_conflicts_detected	Indicates the number of transactions that did not pass the conflict detection check.
Count_transactions_validating	Indicates the current size of the conflict detection database (against which each transaction is certified).
Transactions_committed_all_members	Indicates those transactions that have been successfully committed on all members of the current view. This is updated at a fixed time interval.
Last_conflict_free_transaction	Shows the transaction identifier of the last conflict free transaction checked.

These fields are important for monitoring the performance of the members connected in the group. For example, suppose that one of the group's members is delayed and is not able to keep up to date with the other members of the group. In this case you might see a large number of transactions in the queue. Based on this information, you could decide to either remove the member from the group or delay the processing of transactions on the other members of the group thus reducing the number of queued transactions. This information can can also help you to decide how to adjust the flow control of the Group Replication plugin.

21.3.2 Replication_group_members

This table is used for monitoring the status of the different server instances that are tracked in the current view, or in other words are part of the group and as such are tracked by the membership service.

Table 21.2 replication_group_members

Field	Description
Channel_name	The name of the Group Replication channel.
Member_id	The member server UUID.

Field	Description
Member_host	The network address of the member.
Member_port	The MySQL connection port on which this member is listening.
Member_state	Gives the status of the members of the group (whether they are `ONLINE`, `ERROR`, `RECOVERING`, `OFFLINE` or `UNREACHABLE`)

For more information about the `Member_host` value and it's impact on the distributed recovery process, see Section 21.2.1.3, “User Credentials”.

21.3.3 Replication_connection_status

When connected to a group, some fields in this table show information regarding Group Replication. For instance, the transactions that have been received from the group and queued in the applier queue (the relay log).

Table 21.3 replication_connection_status

Field	Description
Channel_name	The name of the Group Replication channel.
Group_name	Shows the value of the name of the group. It is always a valid UUID.
Source_UUID	Shows the identifier for the group. It is similar to the group name and it is used as the UUID for all the transactions that are generated during group replication.
Service_state	Shows whether the member is a part of the group or not. The possible values of service state can be {ON, OFF and CONNECTING};
Received_transaction_set	Transactions in this GTID set have been received by this member of the group.

21.3.4 Replication_applier_status

The state of the Group Replication related channels and threads can be observed using the regular replication_applier_status table as well. If there are many different worker threads applying transactions, then the worker tables can also be used to monitor what each worker thread is doing.

Table 21.4 replication_applier_status

Field	Description
Channel_name	The name of the Group Replication channel.
Service_state	Shows whether the applier service is ON or OFF.
Remaining_delay	Shows whether there is some applier delay configured.
Count_transactions_retries	The number of retries performed while applying a transaction.
Received_transaction_set	Transactions in this GTID set have been received by this member of the group.

21.3.5 Group Replication Server States

The table `replication_group_members` is updated whenever there is a view change, for example when the configuration of the group is dynamically changed. At that point, servers exchange some of their metadata to synchronize themselves and continue to cooperate together.

There are various states that a server instance can be in. If servers are communicating properly, all report the same states for all servers. However, if there is a network partition, or a server leaves the group, then different information may be reported, depending on which server is queried. Note that if the server has left the group then obviously it cannot report updated information about other servers' state. If there is a partition, such that quorum is lost, servers are not able to coordinate between themselves. As a consequence, they cannot guess what the status of different servers is. Therefore, instead of guessing their state they report that some servers are unreachable.

Table 21.5 Server State

Field	Description	Group Synchronized
ONLINE	The member is ready to serve as a fully functional group member, meaning that the client can connect and start executing transactions.	Yes
RECOVERING	The member is in the process of becoming an active member of the group and is currently going through the recovery process, receiving state information from a donor.	No
OFFLINE	The plugin is loaded but the member does not belong to any group.	No
ERROR	The state of the member. Whenever there is an error on the recovery phase or while applying changes, the server enters this state.	No
UNREACHABLE	Whenever the local failure detector suspects that a given server is not reachable, because maybe it has crashed or was disconnected involuntarily, it shows that server's state as 'unreachable'.	No

Important

Once an instance enters `ERROR` state, the `super_read_only` option is set to `ON`. To leave the `ERROR` state you must manually configure the instance with `super_read_only=OFF`.

Note that Group Replication is *not* synchronous, but eventually synchronous. More precisely, transactions are delivered to all group members in the same order, but their execution is not synchronized, meaning that after a transaction is accepted to be committed, each member commits at its own pace.

21.4 Group Replication Operations

This section describes the different modes of deploying Group Replication, explains common operations for managing groups and provides information about how to tune your groups .

21.4.1 Deploying in Multi-Primary or Single-Primary Mode

Group Replication operates in the following different modes:

- single-primary mode
- multi-primary mode

The default mode is single-primary. It is not possible to have members of the group deployed in different modes, for example one configured in multi-primary mode while another one is in single-primary mode. To switch between modes, the group and not the server, needs to be restarted with a different operating configuration. Regardless of the deployed mode, Group Replication does not

handle client-side fail-over, that must be handled by the application itself, a connector or a middleware framework such as a proxy or router.

When deployed in multi-primary mode, statements are checked to ensure they are compatible with the mode. The following checks are made when Group Replication is deployed in multi-master mode:

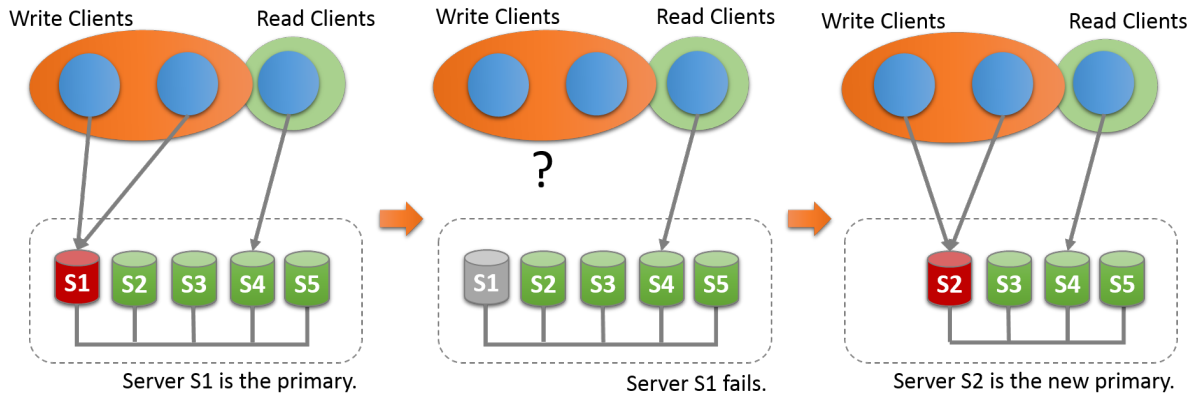
- If a transaction is executed under the SERIALIZABLE isolation level, then its commit fails when synchronizing itself with the group.
- If a transaction executes against a table that has foreign keys with cascading constraints, then the transaction fails to commit when synchronizing itself with the group.

These checks can be deactivated by setting the option `group_replication_enforce_update_everywhere_checks` to `FALSE`. When deploying in single-primary mode, this option *must* be set to `FALSE`.

21.4.1.1 Single-Primary Mode

In this mode the group has a single-primary server that is set to read-write mode. All the other members in the group are set to read-only mode (with `super-read-only=ON`). This happens automatically. The primary is typically the first server to boostrap the group, all other servers that join automatically learn about the primary server and are set to read only.

Figure 21.5 New Primary Election



When in single-primary mode, some of the checks deployed in multi-primary mode are disabled, because the system enforces that only a single writer server is in the group at a time. For example, changes to tables that have cascading foreign keys are allowed, whereas in multi-primary mode they are not. Upon primary member failure, an automatic primary election mechanism chooses the next primary member. The next primary is selected by ordering the remaining servers lexicographically (using their UUID) and picking the first member in the list.

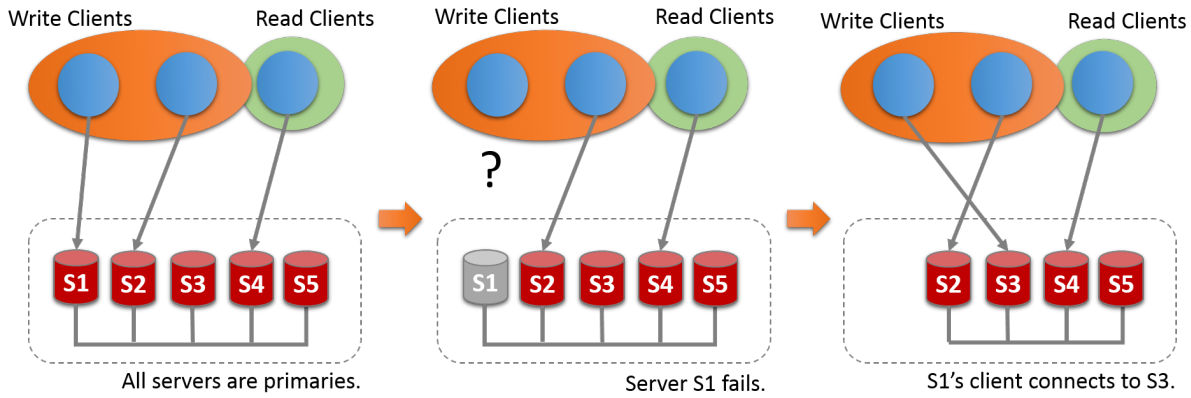
In the event that the primary member leaves the group, then an election is performed and a new primary is chosen from the remaining servers in the group. This election is performed by looking at the new view, ordering the server UUIDs in lexicographical order and by picking the first one. Once a new primary is elected, it is automatically set to read-write and the other secondaries remain as secondaries, and as such, read-only.

It is a good practice to wait for the new primary to apply its replication related relay-log before re-routing the client applications to it.

21.4.1.2 Multi-Primary Mode

In multi-primary mode, there is no notion of a single primary. There is no need to engage an election procedure since there is no server playing any special role.

Figure 21.6 Client Failover



All servers are set to read-write mode when joining the group.

21.4.1.3 Finding the Primary

The following example shows how to find out which server is currently the primary when deployed in single-primary mode.

```
mysql> SELECT VARIABLE_VALUE FROM performance_schema.global_status WHERE VARIABLE_NAME= 'group_replicat
+--------------------------------------+
| VARIABLE_VALUE                       |
+--------------------------------------+
| 69e1a3b8-8397-11e6-8e67-bf68cbc061a4 |
+--------------------------------------+
1 row in set (0,00 sec)
```

21.4.2 Tuning Recovery

Whenever a new member joins a replication group, it connects to a suitable donor and fetches the data that it has missed up until the point it is declared online. This critical component in Group Replication is fault tolerant and configurable. The following section explains how recovery works and how to tune the settings

Donor Selection

A random donor is selected from the existing online members in the group. This way there is a good chance that the same server is not selected more than once when multiple members enter the group.

If the connection to the selected donor fails, a new connection is automatically attempted to a new candidate donor. Once the connection retry limit is reached the recovery procedure terminates with an error.

Note

A donor is picked randomly from the list of online members in the current view.

Enhanced Automatic Donor Switchover

The other main point of concern in recovery as a whole is to make sure that it copes with failures. Hence, Group Replication provides robust error detection mechanisms. In earlier versions of Group Replication, when reaching out to a donor, recovery could only detect connection errors due to authentication issues or some other problem. The reaction to such problematic scenarios was to switch over to a new donor thus a new connection attempt was made to a different member.

This behavior was extended to also cover other failure scenarios:

- *Purged data scenarios* - If the selected donor contains some purged data that is needed for the recovery process then an error occurs. Recovery detects this error and a new donor is selected.
- *Duplicated data* - If a joiner already contains some data that conflicts with the data coming from the selected donor during recovery then an error occurs. This could be caused by some errant transactions present in the joiner.

  One could argue that recovery should fail instead of switching over to another donor, but in heterogeneous groups there is chance that other members share the conflicting transactions and others do not. For that reason, upon error, recovery selects another donor from the group.
- *Other errors* - If any of the recovery threads fail (receiver or applier threads fail) then an error occurs and recovery switches over to a new donor.

Note

In case of some persistent failures or even transient failures recovery automatically retries connecting to the same or a new donor.

Donor Connection Retries

The recovery data transfer relies on the binary log and existing MySQL replication framework, therefore it is possible that some transient errors could cause errors in the receiver or applier threads. In such cases, the donor switch over process has retry functionality, similar to that found in regular replication.

Number of Attempts

The number of attempts a joiner makes when trying to connect to a donor from the pool of donors is 10. This is configured through the `group_replication_recovery_retry_count` plugin variable . The following command sets the maximum number of attempts to connect to a donor to 10.

```
SET GLOBAL group_replication_recovery_retry_count= 10;
```

Note that this accounts for the global number of attempts that the joiner makes connecting to each one of the suitable donors.

Sleep Routines

The `group_replication_recovery_reconnect_interval` plugin variable defines how much time the recovery process should sleep between donor connection attempts. This variable has its default set to 60 seconds and you can change this value dynamically. The following command sets the recovery donor connection retry interval to 120 seconds.

```
SET GLOBAL group_replication_recovery_reconnect_interval= 120;
```

Note, however, that recovery does not sleep after every donor connection attempt. As the joiner is connecting to different servers and not to the same one over and over again, it can assume that the problem that affects server A maybe not affect server B. As such, recovery suspends only when it has gone through all the possible donors. Once the joiner has tried to connect to all the suitable donors in the group and none remains, the recovery process sleeps for the number of seconds configured by the `group_replication_recovery_reconnect_interval` variable.

21.4.3 Network Partitioning

The group needs to achieve consensus whenever a change that needs to be replicated happens. This is the case for regular transactions but is also required for group membership changes and

some internal messaging that keeps the group consistent. Consensus requires a majority of group members to agree on a given decision. When a majority of group members is lost, the group is unable to progress and blocks because it cannot secure majority or quorum.

Quorum may be lost when there are multiple involuntary failures, causing a majority of servers to be removed abruptly from the group. For example in a group of 5 servers, if 3 of them become silent at once, the majority is compromised and thus no quorum can be achieved. In fact, the remaining two are not able to tell if the other 3 servers have crashed or whether a network partition has isolated these 2 alone and therefore the group cannot be reconfigured automatically.

On the other hand, if servers exit the group voluntarily, they instruct the group that it should reconfigure itself. In practice, this means that a server that is leaving tells others that it is going away. This means that other members can reconfigure the group properly, the consistency of the membership is maintained and the majority is recalculated. For example, in the above scenario of 5 servers where 3 leave at once, if the 3 leaving servers warn the group that they are leaving, one by one, then the membership is able to adjust itself from 5 to 2, and at the same time, securing quorum while that happens.

Note

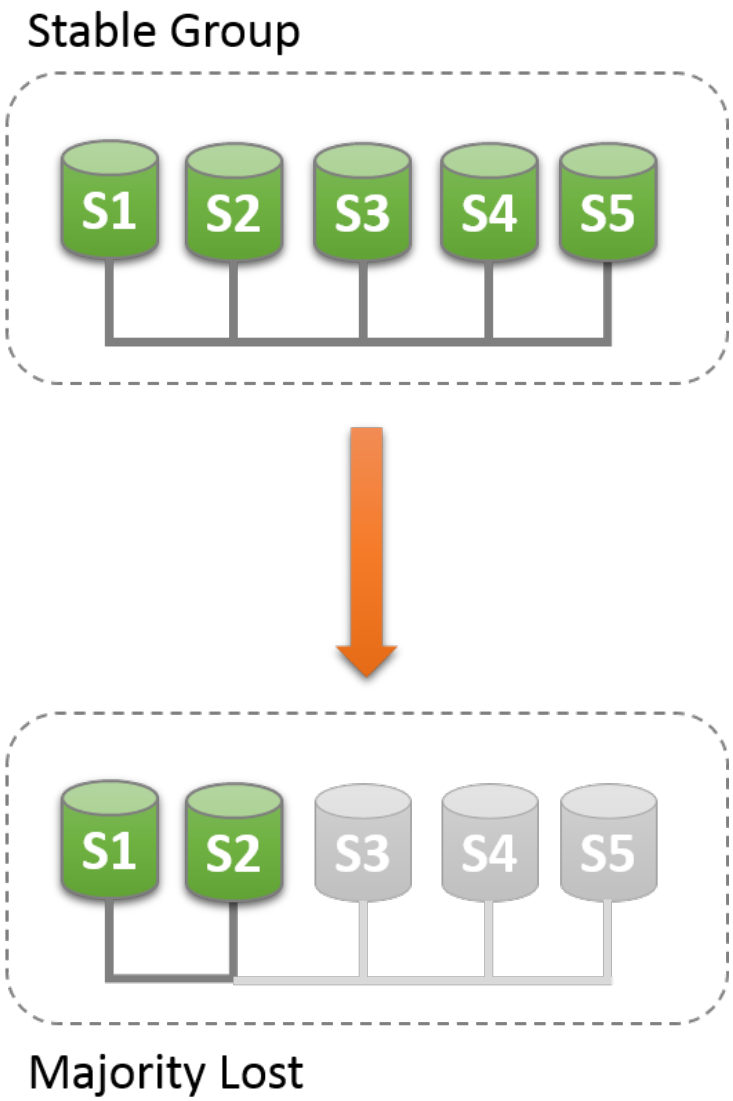
Loss of quorum is by itself a side-effect of bad planning. Plan the group size for the number of expected failures (regardless whether they are consecutive, happen all at once or are sporadic).

The following sections explain what to do if the system partitions in such a way that no quorum is automatically achieved by the servers in the group.

21.4.3.1 Detecting Partitions

The `replication_group_members` performance schema table presents the status of each server in the current view from the perspective of this server. The majority of the time the system does not run into partitioning, and therefore the table shows information that is consistent across all servers in the group. In other words, the status of each server on this table is agreed by all in the current view. However, if there is network partitioning, and quorum is lost, then the table shows the status `UNREACHABLE` for those servers that it cannot contact. This information is exported by the local failure detector built into Group Replication.

Figure 21.7 Losing Quorum



To understand this type of network partition the following section describes a scenario where there are initially 5 servers working together correctly, and the changes that then happen to the group once only 2 servers are online. The scenario is depicted in the figure.

As such, lets assume that there is a group with these 5 servers in it:

- Server s1 with member identifier `199b2df7-4aaf-11e6-bb16-28b2bd168d07`
- Server s2 with member identifier `199bb88e-4aaf-11e6-babe-28b2bd168d07`
- Server s3 with member identifier `1999b9fb-4aaf-11e6-bb54-28b2bd168d07`
- Server s4 with member identifier `19ab72fc-4aaf-11e6-bb51-28b2bd168d07`
- Server s5 with member identifier `19b33846-4aaf-11e6-ba81-28b2bd168d07`

Initially the group is running fine and the servers are happily communicating with each other. You can verify this by logging into s1 and looking at its `replication_group_members` performance schema table. For example:

```
mysql> SELECT * FROM performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+---------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | 1999b9fb-4aaf-11e6-bb54-28b2bd168d07 | 127.0.0.1   |       13002 | ONLINE
| group_replication_applier | 199b2df7-4aaf-11e6-bb16-28b2bd168d07 | 127.0.0.1   |       13001 | ONLINE
| group_replication_applier | 199bb88e-4aaf-11e6-babe-28b2bd168d07 | 127.0.0.1   |       13000 | ONLINE
| group_replication_applier | 19ab72fc-4aaf-11e6-bb51-28b2bd168d07 | 127.0.0.1   |       13003 | ONLINE
| group_replication_applier | 19b33846-4aaf-11e6-ba81-28b2bd168d07 | 127.0.0.1   |       13004 | ONLINE
+---------------------------+--------------------------------------+-------------+-------------+-------
5 rows in set (0,00 sec)
```

However, moments later there is a catastrophic failure and servers s3, s4 and s5 stop unexpectedly. A few seconds after this, looking again at the `replication_group_members` table on s1 shows that it is still online, but several others members are not. In fact, as seen below they are marked as `UNREACHABLE`. Moreover, the system could not reconfigure itself to change the membership, because the majority has been lost.

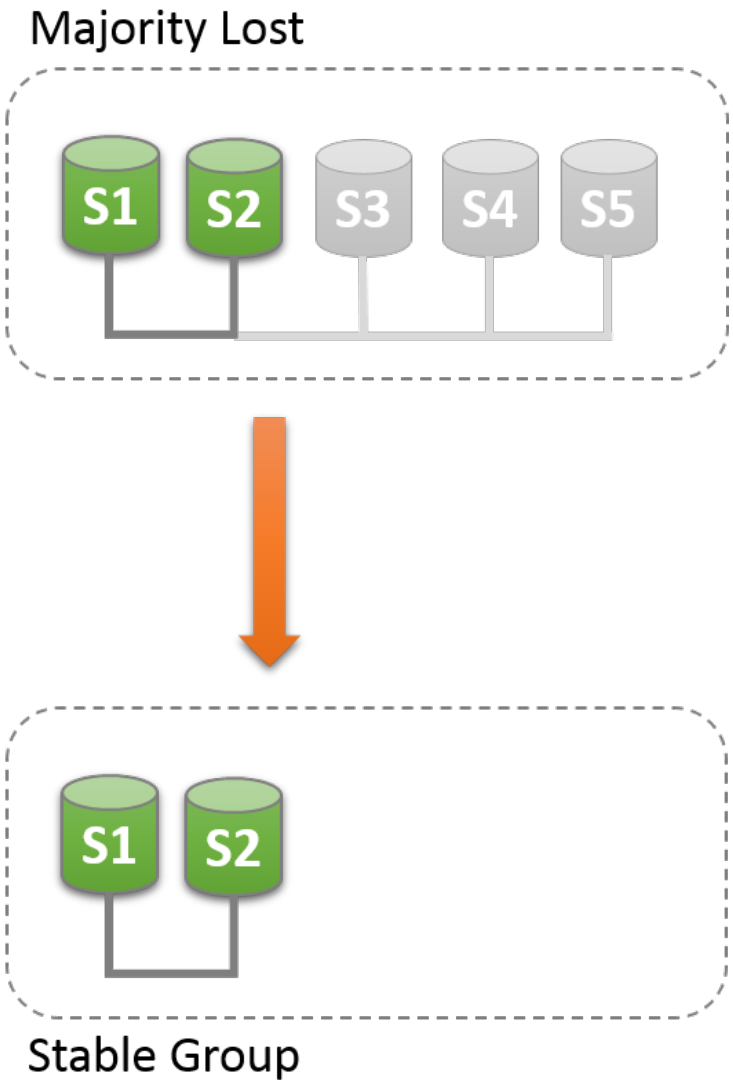
```
mysql> SELECT * FROM performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+---------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | 1999b9fb-4aaf-11e6-bb54-28b2bd168d07 | 127.0.0.1   |       13002 | UNREAC
| group_replication_applier | 199b2df7-4aaf-11e6-bb16-28b2bd168d07 | 127.0.0.1   |       13001 | ONLINE
| group_replication_applier | 199bb88e-4aaf-11e6-babe-28b2bd168d07 | 127.0.0.1   |       13000 | ONLINE
| group_replication_applier | 19ab72fc-4aaf-11e6-bb51-28b2bd168d07 | 127.0.0.1   |       13003 | UNREAC
| group_replication_applier | 19b33846-4aaf-11e6-ba81-28b2bd168d07 | 127.0.0.1   |       13004 | UNREAC
+---------------------------+--------------------------------------+-------------+-------------+-------
5 rows in set (0,00 sec)
```

The table shows that s1 is now in a group that has no means of progressing without external intervention, because a majority of the servers are unreachable. In this particular case, the group membership list needs to be reset to allow the system to proceed, which is explained in this section. Alternatively, you could also choose to stop Group Replication on s1 and s2 (or stop completely s1 and s2), figure out what happened with s3, s4 and s5 and then restart Group Replication (or the servers).

21.4.3.2 Unblocking a Partition

Group replication enables you to reset the group membership list by forcing a specific configuration. For instance in the case above, where s1 and s2 are the only servers online, you could chose to force a membership configuration consisting of only s1 and s2. This requires checking some information about s1 and s2 and then using the `group_replication_force_members` variable.

Figure 21.8 Forcing a New Membership



Suppose that you are back in the situation where s1 and s2 are the only servers left in the group. Servers s3, s4 and s5 have left the group unexpectedly. To make servers s1 and s2 continue, you want to force a membership configuration that contains only s1 and s2.

Warning

This procedure uses `group_replication_force_members` and should be considered a last resort remedy. It *must* be used with extreme care and only for overriding loss of quorum. If misused, it could create an artificial split-brain scenario or block the entire system altogether.

Recall that the system is blocked and the current configuration is the following (as perceived by the local failure detector on s1):

```
mysql> SELECT * FROM performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-----------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER_STA
+---------------------------+--------------------------------------+-------------+-------------+-----------
| group_replication_applier | 1999b9fb-4aaf-11e6-bb54-28b2bd168d07 | 127.0.0.1   |       13002 | UNREACHABL
| group_replication_applier | 199b2df7-4aaf-11e6-bb16-28b2bd168d07 | 127.0.0.1   |       13001 | ONLINE
| group_replication_applier | 199bb88e-4aaf-11e6-babe-28b2bd168d07 | 127.0.0.1   |       13000 | ONLINE
```

```
| group_replication_applier | 19ab72fc-4aaf-11e6-bb51-28b2bd168d07 | 127.0.0.1   |       13003 | UNREAC
| group_replication_applier | 19b33846-4aaf-11e6-ba81-28b2bd168d07 | 127.0.0.1   |       13004 | UNREAC
+---------------------------+--------------------------------------+-------------+-------------+-------
5 rows in set (0,00 sec)
```

The first thing to do is to check what is the peer address (group communication identifier) for s1 and s2. Log in to s1 and s2 and get that information as follows.

```
mysql> SELECT @@group_replication_local_address;
+-----------------------------------+
| @@group_replication_local_address |
+-----------------------------------+
| 127.0.0.1:10000                   |
+-----------------------------------+
1 row in set (0,00 sec)
```

Then log in to s2 and do the same thing:

```
mysql> SELECT @@group_replication_local_address;
+-----------------------------------+
| @@group_replication_local_address |
+-----------------------------------+
| 127.0.0.1:10001                   |
+-----------------------------------+
1 row in set (0,00 sec)
```

Once you know the group communication addresses of s1 (`127.0.0.1:10000`) and s2 (`127.0.0.1:10001`), you can use that on one of the two servers to inject a new membership configuration, thus overriding the existing one that has lost quorum. To do that on s1:

```
mysql> SET GLOBAL group_replication_force_members="127.0.0.1:10000,127.0.0.1:10001";
Query OK, 0 rows affected (7,13 sec)
```

This unblocks the group by forcing a different configuration. Check `replication_group_members` on both s1 and s2 to verify the group membership after this change. First on s1.

```
mysql> select * from performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+---------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | b5ffe505-4ab6-11e6-b04b-28b2bd168d07 | 127.0.0.1   |       13000 | ONLINE
| group_replication_applier | b60907e7-4ab6-11e6-afb7-28b2bd168d07 | 127.0.0.1   |       13001 | ONLINE
+---------------------------+--------------------------------------+-------------+-------------+-------
2 rows in set (0,00 sec)
```

And then on s2.

```
mysql> select * from performance_schema.replication_group_members;
+---------------------------+--------------------------------------+-------------+-------------+-------
| CHANNEL_NAME              | MEMBER_ID                            | MEMBER_HOST | MEMBER_PORT | MEMBER
+---------------------------+--------------------------------------+-------------+-------------+-------
| group_replication_applier | b5ffe505-4ab6-11e6-b04b-28b2bd168d07 | 127.0.0.1   |       13000 | ONLINE
| group_replication_applier | b60907e7-4ab6-11e6-afb7-28b2bd168d07 | 127.0.0.1   |       13001 | ONLINE
+---------------------------+--------------------------------------+-------------+-------------+-------
2 rows in set (0,00 sec)
```

When forcing a new membership configuration, make sure that whatever servers are going to be forced out of the group are indeed stopped. In the scenario depicted above, if s3, s4 and s5 are not really unreachable but instead are online, they may have formed their own functional partition (they are 3 out of 5, hence they have the majority). In that case, forcing a group membership list with s1 and s2 could create an artificial split-brain situation. Therefore it is important before forcing a new membership configuration to ensure that the servers to be excluded are indeed shutdown and if they are not, shut them down before proceding.

21.5 Group Replication Security

This section explains how to secure a group, securing the connections between members of a group, or by establishing a security perimeter using address whitelisting.

21.5.1 IP Address Whitelisting

The Group Replication plugin has a configuration option to determine from which hosts an incoming Group Communication connection can be accepted. This option is called `group_replication_ip_whitelist`. If you set this option on a server s1, then when server s2 is establishing a connection to s1 for the purpose of engaging group communication, then s1 first checks the whitelist before accepting the connection from s2. If s2 is in the whitelist, then s1 accepts the connection, otherwise s1 rejects the connection attempt by s2.

Note

By default, if not specified explicitly, the whitelist is automatically set to the private network addresses that the server has network interfaces on.

If you do not configure any whitelist, the server automatically sets the whitelist to the private networks that the server has an interface on. This means that a server, even if it has interfaces on public IPs, does not by default allow connections from external hosts.

Whenever the IP whitelist is set to AUTOMATIC, an entry in the error log can be emitted in such case, similar to:

```
2016-07-07T06:40:49.320686Z 4 [Note] Plugin group_replication reported: 'Added automatically \\
IP ranges 10.120.40.237/18,10.178.59.44/22,127.0.0.1/8 to the whitelist'
```

You can improve the security of the group further by manually setting the list of IP addresses permitted for group communication connections to come from. The list can be specified in CIDR notation or as simple IP addresses. A comma must separate each entry. For example:

```
mysql> STOP GROUP_REPLICATION;
mysql> SET GLOBAL group_replication_ip_whitelist="10.120.40.237/18,10.178.59.44/22,127.0.0.1/8";
mysql> START GROUP_REPLICATION;
```

Note

The localhost IP address (127.0.0.1) is always added to the whitelist. If not explicitly, it is implicitly and automatically added.

21.5.2 Secure Socket Layer Support (SSL)

MySQL Group Replication supports both OpenSSL and YaSSL builds of MySQL Server.

Group communication connections as well as recovery connections, are secured using SSL. The following sections explain how to configure connections.

Configuring SSL for Recovery

Recovery is performed through a regular asynchronous replication connection. Once the donor is selected, the joiner establishes an asynchronous replication connection. This is all automatic.

However, a user that requires an SSL connection must have been created before the joiner connects to the donor. Typically, this is set up at the time one is provisioning a server to join the group.

```
donor> SET SQL_LOG_BIN=0;
donor> CREATE USER 'rec_ssl_user'@'%' REQUIRE SSL;
donor> GRANT replication slave ON *.* TO 'rec_ssl_user'@'%';
donor> SET SQL_LOG_BIN=1;
```

Assuming that all servers already in the group have a replication user set up to use SSL, you configure the joiner to use those credentials when connecting to the donor. That is done according to the values of the SSL options provided for the Group Replication plugin.

```
new_member> SET GLOBAL group_replication_recovery_use_ssl=1;
new_member> SET GLOBAL group_replication_recovery_ssl_ca= '.../cacert.pem';
new_member> SET GLOBAL group_replication_recovery_ssl_cert= '.../client-cert.pem';
new_member> SET GLOBAL group_replication_recovery_ssl_key= '.../client-key.pem';
```

And by configuring the recovery channel to use the credentials of the user that requires an SSL connection.

```
new_member> CHANGE MASTER TO MASTER_USER="rec_ssl_user" FOR CHANNEL "group_replication_recovery";
new_member> START GROUP_REPLICATION;
```

Configuring SSL for Group Communication

Secure sockets can be used to establish communication between members in a group. The configuration for this depends on the server's SSL configuration. As such, if the server has SSL configured, the Group Replication plugin also has SSL configured. For more information on the options for configuring the server SSL, see Section 6.4.5, “Command Options for Secure Connections”. The options which configure Group Replication are shown in the following table.

Table 21.6 SSL Options

Server Configuration	Plugin Configuration Description
ssl_key	Path of key file. To be used as client and server certificate.
ssl_cert	Path of certificate file. To be used as client and server certificate.
ssl_ca	Path of file with SSL CAs that are trusted.
ssl_capath	Path of directory containing certificates for SSL CAs that are trusted.
ssl_crl	Path of file containing the certificate revocation lists.
ssl_crlpath	Path of directory containing revoked certificate lists files.
ssl_cipher	Permitted ciphers to use while encrypting data over the connection.
tls_version	Secure communication will use this version and its protocols.

These options are MySQL Server configuration options which Group Replication relies on for its configuration. In addition there is the following Group Replication specific option to configure SSL on the plugin itself.

- `group_replication_ssl_mode` - specifies the security state of the connection between Group Replication members.

Table 21.7 group_replication_ssl_mode configuration values

Value	Description
DISABLED	Establish an unencrypted connection (*default*).
REQUIRED	Establish a secure connection if the server supports secure connections.
VERIFY_CA	Like REQUIRED, but additionally verify the server TLS certificate against the configured Certificate Authority (CA) certificates.

Value	Description
VERIFY_IDENTITY	Like VERIFY_CA, but additionally verify that the server certificate matches the host to which the connection is attempted.

The following example shows an example my.cnf file section used to configure SSL on a server and how activate it for Group Replication.

```
[mysqld]
ssl_ca = "cacert.pem"
ssl_capath = "/.../ca_directory"
ssl_cert = "server-cert.pem"
ssl_cipher = "DHE-RSA-AES256-SHA"
ssl_crl = "crl-server-revoked.crl"
ssl_crlpath = "/.../crl_directory"
ssl_key = "server-key.pem"
group_replication_ssl_mode= REQUIRED
```

The only plugin specific configuration option that is listed is `group_replication_ssl_mode`. This option activates the SSL communication between members of the group, by configuring the SSL framework with the `ssl_*` parameters that are provided to the server.

21.5.3 Virtual Private Networks (VPN)

There is nothing preventing Group Replication from operating over a virtual private network. At its core, it just relies on an IPv4 socket to establish connections between servers for the purpose of propagating messages between them.

21.6 Group Replication System Variables

These are the system variables that are specific to the Group Replication plugin. Every configuration option is prefixed with "group_replication".

Important

Although most variables are described as dynamic and can be changed while the server is running, most changes only take effect upon restarting the Group Replication plugin. Variables which can be changed without requiring a restart of the plugin are specifically noted as such in this section.

- `group_replication_group_name`

Command-Line Format	`--group-replication-group-name=value`	
System Variable	Name	`group_replication_group_name`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

The name of the group which this server instance belongs to. Must be a valid UUID.

- `group_replication_start_on_boot`

Command-Line Format	`--group-replication-start-on-boot=value`	
System Variable	Name	`group_replication_start_on_boot`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`ON`

Whether the server should start Group Replication or not during server start.

- `group_replication_local_address`

Command-Line Format	`--group-replication-local-address=value`	
System Variable	**Name**	`group_replication_local_address`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The local address as a `host:port` formatted string.

- `group_replication_member_weight`

Introduced	8.0.2	
Command-Line Format	`--group-replication-member-weight=value`	
System Variable	**Name**	`group_replication_member_weight`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`50`
	Min Value	`0`
	Max Value	`100`

A percentage weight that can be assigned to members to influence the chance of the member being elected as primary in the event that the existing primary leaves the group. Assign numeric weights to members to ensure that specific members are elected, for example during scheduled maintenance of the primary or to ensure certain hardware is prioritised in the event of failover.

- `group_replication_group_seeds`

Command-Line Format	`--group-replication-group-seeds=value`	
System Variable	**Name**	`group_replication_group_seeds`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

A list of peer addresses as a comma separated list such as `host1:port1,host2:port2`. Each address is validated when starting Group Replication and the list must contain at least one valid address. Any invalid host names could cause `START GROUP_REPLICATION` to fail.

- `group_replication_force_members`

Command-Line Format	`--group-replication-force-members=value`	
System Variable	Name	`group_replication_force_members`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

A list of peer addresses as a comma separated list such as `host1:port1,host2:port2`. This option is used to force a new group membership, in which the excluded members do not receive a new view and are blocked. You need to manually kill the excluded servers. Any invalid host names in the list could cause subsequent `START GROUP_REPLICATION` statements to fail because they could block group membership.

- `group_replication_bootstrap_group`

Command-Line Format	`--group-replication-bootstrap-group=value`	
System Variable	Name	`group_replication_bootstrap_group`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Configure this server to bootstrap the group. This option must only be set on one server and only when starting the group for the first time or restarting the entire group. After the group has been bootstrapped, set this option to `OFF`. It should be set to `OFF` both dynamically and in the configuration files. Starting two servers or restarting one server with this option set while the group is running may lead to an artificial split brain situation, where two independent groups with the same name are bootstrapped.

- `group_replication_poll_spin_loops`

Command-Line Format	`--group-replication-poll-spin-loops=value`	
System Variable	Name	`group_replication_poll_spin_loops`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	`0`
	Min Value	`0`
	Max Value	`4294967295`

Permitted Values (64-bit platforms)	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`18446744073709547520`

The number of times the group communication thread waits for the communication engine mutex to be released before the thread waits for more incoming network messages.

- `group_replication_recovery_retry_count`

Command-Line Format	`--group-replication-recovery-retry-count=value`	
System Variable	**Name**	`group_replication_recovery_retry_count`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`10`
	Min Value	`0`
	Max Value	`31536000`

The number of times that the member that is joining tries to connect to the available donors before giving up.

- `group_replication_recovery_reconnect_interval`

Command-Line Format	`--group-replication-recovery-reconnect-interval=value`	
System Variable	**Name**	`group_replication_recovery_reconnect_interval`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`60`
	Min Value	`0`
	Max Value	`0`

The sleep time, in seconds, between reconnection attempts when no donor was found in the group.

- `group_replication_recovery_use_ssl`

Command-Line Format	`--group-replication-recovery-use-ssl=value`	
System Variable	**Name**	`group_replication_recovery_use_ssl`
	Variable Scope	Global

	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Whether Group Replication recovery connection should use SSL or not.

- `group_replication_recovery_ssl_ca`

Command-Line Format	`--group-replication-recovery-ssl-ca=value`	
System Variable	**Name**	`group_replication_recovery_ssl_ca`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The path to a file that contains a list of trusted SSL certificate authorities.

- `group_replication_recovery_ssl_capath`

Command-Line Format	`--group-replication-recovery-ssl-capath=value`	
System Variable	**Name**	`group_replication_recovery_ssl_capath`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The path to a directory that contains trusted SSL certificate authority certificates.

- `group_replication_recovery_ssl_cert`

Command-Line Format	`--group-replication-recovery-ssl-cert=value`	
System Variable	**Name**	`group_replication_recovery_ssl_cert`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The name of the SSL certificate file to use for establishing a secure connection.

- `group_replication_recovery_ssl_key`

Command-Line Format	`--group-replication-recovery-ssl-key=value`	
System Variable	**Name**	`group_replication_recovery_ssl_key`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string

The name of the SSL key file to use for establishing a secure connection.

- `group_replication_recovery_ssl_cipher`

Command-Line Format	`--group-replication-recovery-ssl-cipher=value`	
System Variable	Name	`group_replication_recovery_ssl_cipher`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

A list of permissible ciphers to use for SSL encryption.

- `group_replication_recovery_ssl_crl`

Command-Line Format	`--group-replication-recovery-ssl-crl=value`	
System Variable	Name	`group_replication_recovery_ssl_crl`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

The path to a directory that contains files containing certificate revocation lists.

- `group_replication_recovery_ssl_crlpath`

Command-Line Format	`--group-replication-recovery-ssl-crlpath=value`	
System Variable	Name	`group_replication_recovery_ssl_crlpath`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	string

The path to a directory that contains files containing certificate revocation lists.

- `group_replication_recovery_ssl_verify_server_cert`

Command-Line Format	`--group-replication-recovery-ssl-verify-server-cert=value`	
System Variable	Name	`group_replication_recovery_ssl_verify_server_cert`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Make the recovery process check the server's Common Name value in the donor sent certificate.

- `group_replication_recovery_complete_at`

Command-Line Format	`--group-replication-recovery-complete-at=value`	
System Variable	Name	`group_replication_recovery_complete_at`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	`TRANSACTIONS_APPLIED`
	Valid Values	`TRANSACTIONS_CERTIFIED`
		`TRANSACTIONS_APPLIED`

Recovery policies when handling cached transactions after state transfer. This option specifies whether a member is marked online after it has received all transactions that it missed before it joined the group (`TRANSACTIONS_CERTIFIED`) or after it has received and applied them (`TRANSACTIONS_APPLIED`).

- `group_replication_components_stop_timeout`

Command-Line Format	`--group-replication-components-stop-timeout=value`	
System Variable	Name	`group_replication_components_stop_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	`31536000`

Timeout, in seconds, that Group Replication waits for each of the components when shutting down.

- `group_replication_allow_local_lower_version_join`

Command-Line Format	`--group-replication-allow-local-lower-version-join=value`	
System Variable	Name	`group_replication_allow_local_lower_version_join`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean
	Default	`OFF`

Allow the current server to join the group even if it has a lower plugin version than the group.

- `group_replication_allow_local_disjoint_gtids_join`

Command-Line Format	`--group-replication-allow-local-disjoint-gtids-join=value`	
System Variable	Name	`group_replication_allow_local_disjoint_gtids_join`

	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Allow the current server to join the group even if it has transactions not present in the group.

Warning

Use caution when enabling this option as incorrect usage could lead to inconsistencies in the group.

- `group_replication_auto_increment_increment`

Command-Line Format	`--group-replication-auto-increment-increment=value`	
System Variable	**Name**	`group_replication_auto_increment_increment`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`7`
	Min Value	`1`
	Max Value	`65535`

Determines the interval between successive column values for transactions that execute on this server instance.

- `group_replication_compression_threshold`

Command-Line Format	`--group-replication-compression-threshold=value`	
System Variable	**Name**	`group_replication_compression_threshold`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1000000`
	Min Value	`0`
	Max Value	`4294967296`

The value in bytes above which (LZ4) compression is enforced. When set to zero, deactivates compression.

- `group_replication_gtid_assignment_block_size`

Command-Line Format	--group-replication-gtid-assignment-block-size=value	
System Variable	Name	group_replication_gtid_assignment_block_size
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values (32-bit platforms)	Type	integer
	Default	1000000
	Min Value	1
	Max Value	4294967295
Permitted Values (64-bit platforms)	Type	integer
	Default	1000000
	Min Value	1
	Max Value	18446744073709547520

The number of consecutive GTIDs that are reserved for each member. Each member consumes its blocks and reserves more when needed.

- group_replication_ssl_mode

Command-Line Format	--group-replication-ssl-mode=value	
System Variable	Name	group_replication_ssl_mode
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	enumeration
	Default	DISABLED
	Valid Values	DISABLED
		REQUIRED
		VERIFY_CA
		VERIFY_IDENTITY

Specifies the security state of the connection between Group Replication members.

- group_replication_single_primary_mode

Command-Line Format	--group-replication-single-primary-mode=value	
System Variable	Name	group_replication_single_primary_mode
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	boolean

	Default	`ON`

Instructs the group to automatically pick a single server to be the one that handles read/write workload. This server is the PRIMARY and all others are SECONDARIES.

- `group_replication_transaction_size_limit`

Command-Line Format	`--group-replication-transaction-size-limit=value`	
System Variable	**Name**	`group_replication_transaction_size_limit`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`2147483647`
Permitted Values (>= 8.0.2)	**Type**	integer
	Default	`150000000`
	Min Value	`0`
	Max Value	`2147483647`

Configures the maximum transaction size in bytes which the group accepts. Transactions larger than this size are rolled back. Use this option to avoid large transactions causing the group to fail. A large transaction can cause problems for a group, either in terms of memory allocation or network bandwidth consumption, which may cause the failure detector to trigger because a given member is unreachable while it is busy processing the large transaction. When set to 0 there is no limit to the size of transactions the group accepts, and there may be the risk of large transactions causing the group to fail. Adjust the value of this variable depending on the size of workload you require from the group.

- `group_replication_unreachable_majority_timeout`

Introduced	8.0.2	
Command-Line Format	`--group-replication-unreachable-majority-timeout=value`	
System Variable	**Name**	`group_replication_unreachable_majority_timeout`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`31536000`

Configures how long members that suffer a network partition and cannot connect to the majority wait before leaving the group. By default set to 0, which means that members that find themselves in a minority due to a network partition wait forever to connect the group. In a group of 5 servers (S1,S2,S3,S4,S5), if there is a disconnection between (S1,S2) and (S3,S4,S5) there is a network partition. The first group (S1,S2) is now in a minority because it cannot contact more than half of the group. While the majority group (S3,S4,S5) remains running, the minority group waits forever for a network reconnection. Any transactions processed by the minority group are blocked until Group Replication is stopped using `STOP GROUP REPLICATION` on the members of the minority.

If configured to a number of seconds, members wait for this amount of time after losing contact with the majority of members before leaving the group. All pending transactions processed by the minority are rolled back and the servers in the minority partition move to the `ERROR` state and set themselves to `super_read_only=ON` mode.

Warning

When you have a symmetric group, with just two members for example (S0,S2), if there is a network partition and there is no majority, after the configured timeout all members shut down and enter `ERROR` state.

- `group_replication_enforce_update_everywhere_checks`

Command-Line Format	`--group-replication-enforce-update-everywhere-checks=value`	
System Variable	**Name**	`group_replication_enforce_update_everywhere_checks`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	boolean
	Default	`OFF`

Enable or disable strict consistency checks for multi-primary update everywhere.

- `group_replication_flow_control_mode`

Command-Line Format	`--group-replication-flow-control-mode=value`	
System Variable	**Name**	`group_replication_flow_control_mode`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	enumeration
	Default	`QUOTA`
	Valid Values	`DISABLED`
		`QUOTA`

Specifies the mode used for flow control. This variable can be changed without resetting Group Replication.

- `group_replication_flow_control_hold_percent`

Introduced	8.0.2	

Command-Line Format	--group-replication-flow-control-hold-percent=value	
System Variable	Name	group_replication_flow_control_hold_percent
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	10
	Min Value	0
	Max Value	100

Defines what percentage of the group quota remains unused to allow a cluster under flow control to catch up on backlog. A value of 0 implies that no part of the quota is reserved for catching up on the work backlog.

- group_replication_flow_control_certifier_threshold

Command-Line Format	--group-replication-flow-control-certifier-threshold=value	
System Variable	Name	group_replication_flow_control_certifier_threshold
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	25000
	Min Value	0
	Max Value	2147483648

Specifies the number of waiting transactions in the certifier queue that trigger flow control. This variable can be changed without resetting Group Replication.

- group_replication_flow_control_applier_threshold

Command-Line Format	--group-replication-flow-control-applier-threshold=value	
System Variable	Name	group_replication_flow_control_applier_threshold
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	Type	integer
	Default	25000
	Min Value	0

	Max Value	2147483648

Specifies the number of waiting transactions in the applier queue that trigger flow control. This variable can be changed without resetting Group Replication.

- `group_replication_flow_control_max_commit_quota`

Introduced	8.0.2	
Command-Line Format	`--group-replication-flow-control-max-commit-quota=value`	
System Variable	**Name**	`group_replication_flow_control_max_commit_quota`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`2147483648`

Defines the maximum flow control quota of the group, or the maximum available quota for any period while flow control is enabled. A value of 0 implies that there is no maximum quota set. Cannot be smaller than `group_replication_flow_control_min_quota` and `group_replication_flow_control_min_recovery_quota`.

- `group_replication_flow_control_member_quota_percent`

Introduced	8.0.2	
Command-Line Format	`--group-replication-flow-control-member-quota-percent=value`	
System Variable	**Name**	`group_replication_flow_control_member_quota_percent`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`100`

Defines the percentage of the quota that a member should assume is available for itself when calculating the quotas. A value of 0 implies that the quota should be split equally between members that were writers in the last period.

- `group_replication_flow_control_min_quota`

Introduced	8.0.2	

Command-Line Format	`--group-replication-flow-control-min-quota=value`	
System Variable	**Name**	`group_replication_flow_control_min_quota`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`2147483648`

Controls the lowest flow control quota that can be assigned to a member, independently of the calculated minimum quota executed in the last period. A value of 0 implies that there is no minimum quota. Cannot be larger than `group_replication_flow_control_max_commit_quota`.

- `group_replication_flow_control_min_recovery_quota`

Introduced	8.0.2	
Command-Line Format	`--group-replication-flow-control-min-recovery-quota=value`	
System Variable	**Name**	`group_replication_flow_control_min_recovery_quota`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`0`
	Min Value	`0`
	Max Value	`2147483648`

Controls the lowest quota that can be assigned to a member because of another recovering member in the group, independently of the calculated minimum quota executed in the last period. A value of 0 implies that there is no minimum quota. Cannot be larger than `group_replication_flow_control_max_commit_quota`.

- `group_replication_flow_control_period`

Introduced	8.0.2	
Command-Line Format	`--group-replication-flow-control-period=value`	
System Variable	**Name**	`group_replication_flow_control_period`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`1`

	Min Value	1
	Max Value	60

Defines how many seconds to wait between flow control iterations, in which flow control messages are sent and flow control management tasks are run.

- `group_replication_flow_control_release_percent`

Introduced	8.0.2	
Command-Line Format	`--group-replication-flow-control-release-percent=value`	
System Variable	**Name**	`group_replication_flow_control_release_percent`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	integer
	Default	`50`
	Min Value	`0`
	Max Value	`1000`

Defines how the group quota should be released when flow control no longer needs to throttle the writer members, with this percentage being the quota increase per flow control period. A value of 0 implies that once the flow control thresholds are within limits the quota is released in a single flow control iteration. The range allows the quota to be released at up to 10 times current quota, as that allows a greater degree of adaptation, mainly when the flow control period is large and the quotas are very small.

- `group_replication_ip_whitelist`

Command-Line Format	`--group-replication-ip-whitelist=value`	
System Variable	**Name**	`group_replication_recovery_ip_whitelist`
	Variable Scope	Global
	Dynamic Variable	Yes
Permitted Values	**Type**	string
	Default	`AUTOMATIC`

Specifies which hosts are permitted to connect to the group. By default set to `AUTOMATIC`, which permits connections from private subnetworks active on the host. Active interfaces on the host are scanned and those with addresses on private subnetworks are automatically added to the permitted list. Alternatively you can specify the permitted hosts using a comma separated list of IPv4 addresses or subnet CIDR notation. For example `192.168.1.0/24,10.0.0.1`.

Note

Address 127.0.0.1 is always permitted to connect, even when not specified using `group_replication_ip_whitelist`.

21.7 Requirements and Limitations

This section lists and explains the requirements and limitations of Group Replication.

21.7.1 Group Replication Requirements

Server instances that you want to use for Group Replication must satisfy the following requirements.

Infrastructure

- **InnoDB Storage Engine.** Data must be stored in the `InnoDB` transactional storage engine. Transactions are executed optimistically and then, at commit time, are checked for conflicts. If there are conflicts, in order to maintain consistency across the group, some transactions are rolled back. This means that a transactional storage engine is required. Moreover, `InnoDB` provides some additional functionality that enables better management and handling of conflicts when operating together with Group Replication.

- **Primary Keys.** Every table that is to be replicated by the group must have an explicit primary key defined. Primary keys play an extremely important role in determining which transactions conflict by identifying exactly which rows each transaction has modified.

- **IPv4 Network.** The group communication engine used by MySQL Group Replication only supports IPv4. Therefore, Group Replication requires an IPv4 network infrastructure.

- **Network Performance.** Group Replication is designed to be deployed in a cluster environment where server instances are very close to each other, and is impacted by both network latency as well as network bandwidth.

Server Instance Configuration

The following options must be configured on server instances that are members of a group.

- **Binary Log Active.** Set `--log-bin[=log_file_name]`. MySQL Group Replication replicates binary log contents, therefore the binary log needs to be turned on for it to operate. See Section 5.4.4, “The Binary Log”.

- **Slave Updates Logged.** Set `--log-slave-updates`. Servers need to log binary logs that are applied through the replication applier. Servers in the group need to log all transactions that they receive and apply from the group. This is required because recovery is conducted by relying on binary logs form participants in the group. Therefore, copies of each transaction need to exist on every server, even for those transactions that were not initiated on the server itself.

- **Binary Log Row Format.** Set `--binlog-format=row`. Group Replication relies on row-based replication format to propagate changes consistently among the servers in the group. It relies on row-based infrastructure to be able to extract the necessary information to detect conflicts among transactions that execute concurrently in different servers in the group. See Section 18.2.1, “Replication Formats”.

- **Global Transaction Identifiers On.** Set `--gtid-mode=ON`. Group Replication uses global transaction identifiers to track exactly which transactions have been committed on every server instance and thus be able to infer which servers have executed transactions that could conflict with already committed transactions elsewhere. In other words, explicit transaction identifiers are a fundamental part of the framework to be able to determine which transactions may conflict. See Section 18.1.3, “Replication with Global Transaction Identifiers”.

- **Replication Information Repositories.** Set `--master-info-repository=TABLE` and `--relay-log-info-repository=TABLE`. The replication applier needs to have the master information and relay log metadata written to the `mysql.slave_master_info` and `mysql.slave_relay_log_info` system tables. This ensures the Group Replication plugin

has consistent recoverability and transactional management of the replication metadata. See Section 18.2.4.2, “Slave Status Logs”.

- **Transaction Write Set Extraction.** Set `--transaction-write-set-extraction=XXHASH64` so that while collecting rows to log them to the binary log, the server collects the write set as well. The write set is based on the primary keys of each row and is a simplified and compact view of a tag that uniquely identifies the row that was changed. This tag is then used for detecting conflicts.

- **Multi-threaded Appliers.** Group Replication members can be configured as multi-threaded appliers, enabling transactions to be applied in parallel. Set `--slave-parallel-workers=N` (where `N` is the number of parallel applier threads), `--slave-preserve-commit-order=1`, and `--slave-parallel-type=LOGICAL_CLOCK`. Setting `--slave-parallel-workers=N` enables the multi-threaded applier on the member. Group Replication relies on consistency mechanisms built around the guarantee that all participating members receive and apply committed transaction in the same order, so you must also set `--slave-preserve-commit-order=1` to ensure that the final commit of parallel transactions is in the same order as the original transactions. Finally, in order to determine which transactions can be executed in parallel, the relay log must contain transaction parent information generated with `--slave-parallel-type=LOGICAL_CLOCK`. Attempting to add a member with `--slave-parallel-workers` set to greater than 0 without also setting the other two options, generates an error and the instance is prevented from joining.

21.7.2 Group Replication Limitations

The following known limitations exist for Group Replication.

- **Replication Event Checksums.** Due to a design limitation of replication event checksums, Group Replication cannot currently make use of them. Therefore set `--binlog-checksum=NONE`.

- **Gap Locks.** The certification process does not take into account gap locks, as information about gap locks is not available outside of `InnoDB`. See Gap Locks for more information.

Note

Unless you rely on `REPEATABLE READ` semantics in your applications, we recommend using the `READ COMMITTED` isolation level with Group Replication. InnoDB does not use gap locks in `READ COMMITTED`, which aligns the local conflict detection within InnoDB with the distributed conflict detection performed by Group Replication.

- **Table Locks and Named Locks.** The certification process does not take into account table locks (see Section 13.3.5, “LOCK TABLES and UNLOCK TABLES Syntax”) or named locks (see `GET_LOCK()`).

- **SERIALIZABLE Isolation Level.** `SERIALIZABLE` isolation level is not supported in multi-primary groups by default. Setting a transaction isolation level to `SERIALIZABLE` configures Group Replication to refuse to commit the transaction.

- **Concurrent DDL versus DML Operations.** Concurrent data definition statements and data manipulation statements executing against the same object but on different servers is not supported when using multi-primary mode. During execution of Data Definition Language (DDL) statements on an object, executing concurrent Data Manipulation Language (DML) on the same object but on a different server instance has the risk of conflicting DDL executing on different instances not being detected.

- **Foreign Keys with Cascading Constraints.** Multi-primary mode groups (members all configured with `group_replication_single_primary_mode=OFF`) do not support tables with multi-level foreign key dependencies, specifically tables that have defined `CASCADING` foreign key constraints. This is because foreign key constraints that result in cascading operations executed by a multi-primary mode group can result in undetected conflicts and

lead to inconsistent data across the members of the group. Therefore we recommend setting `group_replication_enforce_update_everywhere_checks=ON` on server instances used in multi-primary mode groups to avoid undetected conflicts.

In single-primary mode this is not a problem as it does not allow concurrent writes to multiple members of the group and thus there is no risk of undetected conflicts.

- **Very Large Transactions.** Individual transactions that result in GTID contents which are large enough that it cannot be copied between group members over the network within a 5 second window can cause failures in the group communication. To avoid this issue try and limit the size of your transactions as much as possible. For example, split up files used with `LOAD DATA INFILE` into smaller chunks.

21.8 Frequently Asked Questions

This section provides answers to frequently asked questions.

What is the maximum number of MySQL servers in a group?

A group can consist of maximum 9 servers. Attempting to add another server to a group with 9 members causes the request to join to be refused.

How are servers in a group connected?

Servers in a group connect to the other servers in the group by opening a peer-to-peer TCP connection. These connections are only used for internal communication and message passing between servers in the group.

What is the group_replication_bootstrap_group option used for?

The bootstrap flag instructs a member to *create* a group and act as the initial seed server. The second member joining the group needs to ask the member that bootstrapped the group to dynamically change the configuration in order for it to be added to the group.

A member needs to bootstrap the group in two scenarios. When the group is originally created, or when shutting down and restarting the entire group.

How do I set credentials for the recovery procedure?

You pre-configure the Group Replication recovery channel credentials using the `CHANGE MASTER TO` statement.

Can I scale-out my write-load using Group Replication?

Not directly, but MySQL Group replication is a shared nothing full replication solution, where all servers in the group replicate the same amount of data. Therefore if one member in the group writes N bytes to storage as the result of a transaction commit operation, then roughly N bytes are written to storage on other members as well, because the transaction is replicated everywhere.

However, given that other members do not have to do the same amount of processing that the original member had to do when it originally executed the transaction, they apply the changes faster. Transactions are replicated in a format that is used to apply row transformations only, without having to re-execute transactions again (row-based format).

Furthermore, given that changes are propagated and applied in row-based format, this means that they are received in an optimized and compact format, and likely reducing the number of IO operations required when compared to the originating member.

To summarize, you can scale-out processing, by spreading conflict free transactions throughout different members in the group. And you can likely scale-out a small fraction of your IO operations, since remote servers receive only the necessary changes to read-modify-write changes to stable storage.

Does Group Replication require more network bandwidth and CPU, when compared to simple replication and under the same workload?

Some additional load is expected because servers need to be constantly interacting with each other for synchronization purposes. It is difficult to quantify how much more data. It also depends on the size of the group (three servers puts less stress on the bandwidth requirements than nine servers in the group).

Also the memory and CPU footprint are larger, because more complex work is done for the server synchronization part and for the group messaging.

Can I deploy Group Replication across wide-area networks?

Yes, but the network connection between each member *must* be reliable and have suitable perfomance. Low latency, high bandwidth network connections are a requirement for optimal performance.

If network bandwidth alone is an issue, then Section 21.9.7.2, “Message Compression” can be used to lower the bandwidth required. However, if the network drops packets, leading to re-transmissions and higher end-to-end latency, throughput and latency are both negatively affected.

Warning

When the network round-trip time (RTT) between any group members is 2 seconds or more you could encounter problems as the built-in failure detection mechanism could be incorrectly triggered.

Do members automatically rejoin a group in case of temporary connectivity problems?

This depends on the reason for the connectivity problem. If the connectivity problem is transient and the reconnection is quick enough that the failure detector is not aware of it, then the server may not be removed from the group. If it is a "long" connectivity problem, then the failure detector eventually suspects a problem and the server is removed from the group.

Once a server is removed from the group, you need to join it back again. In other words, after a server is removed explicitly from the group you need to rejoin it manually (or have a script doing it automatically).

When is a member excluded from a group?

If the member becomes silent, the other members remove it from the group configuration. In practice this may happen when the member has crashed or there is a network disconnection.

The failure is detected after a given timeout elapses for a given member and a new configuration without the silent member in it is created.

What happens when one node is significantly lagging behind?

There is no method for defining policies for when to expel members automatically from the group. You need to find out why a member is lagging behind and fix that or remove the member from the group.

Otherwise, if the server is so slow that it triggers the flow control, then the entire group slows down as well. The flow control can be configured according to the your needs.

Upon suspicion of a problem in the group, is there a special member responsible for triggering a reconfiguration?

No, there is no special member in the group in charge of triggering a reconfiguration.

Any member can suspect that there is a problem. All members need to (automatically) agree that a given member has failed. One member is in charge of expelling it from the group, by triggering a reconfiguration. Which member is responsible for expelling the member is not something you can control or set.

Can I use Group Replication for sharding?

Group Replication is designed to provide highly available replica sets; data and writes are duplicated on each member in the group. For scaling beyond what a single system can provide, you need an orchestration and sharding framework built around a number of Group Replication sets, where each replica set maintains and manages a given shard or partition of your total dataset. This type of setup, often called a “sharded cluster”, allows you to scale reads and writes linearly and without limit.

How do I use Group Replication with SELinux?

If SELinux is enabled, which you can verify using `sestatus -v`, then you need to enable the use of the Group Replication communication port, configured by `group_replication_local_address`, for `mysqld` so that it can bind to it and listen there. To see which ports MySQL is currently allowed to use, issue `semanage port -l | grep mysqld`. Assuming the port configured is 6606, add the necessary port to those permitted by SELinux by issuing `semanage port -a -t mysqld_port_t -p tcp 6606`.

How do I use Group Replication with iptables?

If iptables is enabled, then you need to open up the Group Replication port for communication between the machines. To see the current rules in place on each machine, issue `iptables -L`. Assuming the port configured is 6606, enable communication over the necessary port by issuing `iptables -A INPUT -p tcp --dport 6606 -j ACCEPT`.

How do I recover the relay log for a replication channel used by a group member?

The replication channels used by Group Replication behave in the same way as replication channels used in master to slave replication, and as such rely on the relay log. In the event of a change of the `relay_log` variable, or when the option is not set and the host name changes, there is a chance of errors. See Section 18.2.4.1, “The Slave Relay Log” for a recovery procedure in this situation. Alternatively, another way of fixing the issue specifically in Group Replication is to issue a `STOP GROUP_REPLICATION` statement and then a `START GROUP_REPLICATION` statement to restart the instance. The Group Replication plugin creates the `group_replication_applier` channel again.

21.9 Group Replication Technical Details

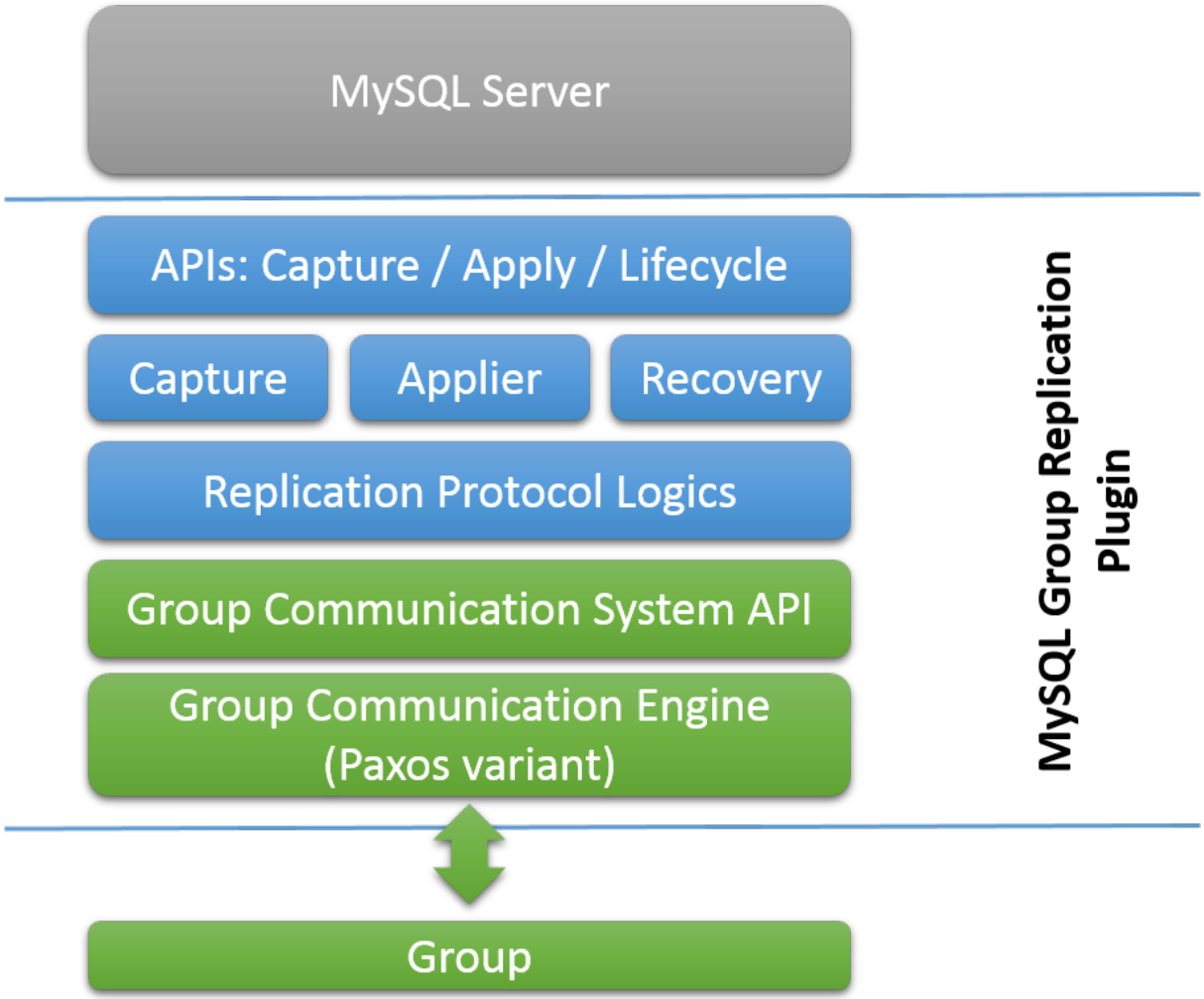
This section provides more technical details about MySQL Group Replication.

21.9.1 Group Replication Plugin Architecture

MySQL Group Replication is a MySQL plugin and it builds on the existing MySQL replication infrastructure, taking advantage of features such as the binary log, row-based logging, and global

transaction identifiers. It integrates with current MySQL frameworks, such as the performance schema or plugin and service infrastructures. The following figure presents a block diagram depicting the overall architecture of MySQL Group Replication.

Figure 21.9 Group Replication Plugin Block Diagram



Starting from the top of the figure, there is a set of APIs that control how the Group Replication plugin interacts with MySQL Server (the gray box). There are interfaces to make information flow from the server to the plugin and vice versa. Such interfaces isolate the server core from the plugin and are mostly hooks placed in the transaction execution pipeline. In one direction, from server to the plugin, there are notifications such as the server is starting, the server is recovering, the server is ready to accept connections, the server is about to commit a transaction, and so on. In the other direction, the plugin instructs the server to commit or abort ongoing transactions, queue transactions in the relay-log, and so on.

Below the API block is a set of components that react when a notification is routed to them. The capture component is responsible for keeping track of context related to transactions that are executing. The applier component is responsible for executing remote transactions on the database. The recovery component manages distributed recovery and is responsible for getting a server that is joining the group up to date by selecting the donor, orchestrating the catch up procedure and reacting on donor failures.

Continuing down the stack, the replication protocol module contains the specific logic of the replication protocol. It handles conflict detection, receives and propagates transactions to the group.

The first green box is the Group Communication API. It is a high level API abstracting the properties required to build a replicated state machine (details in the Background chapter). Therefore it decouples the implementation of the messaging layer from the remaining upper layers of the plugin. The MySQL Group Replication plugin contains an implementation of a Paxos based group communication engine.

21.9.2 The Group

In MySQL Group Replication, a set of servers forms a replication group. A group has a name, which takes the form of a UUID. The group is dynamic and servers can leave (either voluntarily or involuntarily) and join it at any time. The group adjusts itself whenever servers join or leave.

If a server joins the group, it automatically brings itself up to date by fetching the missing state from an existing server. This state is transferred by means of Asynchronous MySQL replication. If a server leaves the group, for instance it was taken down for maintenance, the remaining servers notice that it has left and reconfigure the group automatically. The group membership service described at Section 21.1.3.2, “Group Membership” powers all of this.

21.9.3 Data Manipulation Statements

As there are no primary servers (masters) for any particular data set, every server in the group is allowed to execute transactions at any time, even transactions that change state (RW transactions).

Any server may execute a transaction without any *a priori* coordination. But, at commit time, it coordinates with the rest of the servers in the group to reach a decision on the fate of that transaction. This coordination serves two purposes: (i) check whether the transaction should commit or not; (ii) and propagate the changes so that other servers can apply the transaction as well.

As a transaction is sent through an atomic broadcast, either all servers in the group receive the transaction or none do. If they receive it, then they all receive it in the same order with respect to other transactions that were sent before. Conflict detection is carried out by inspecting and comparing write sets of transactions. Thus, they are detected at the row level. Conflict resolution follows the first committer wins rule. If t1 and t2 execute concurrently at different sites, because t2 is ordered before t1, and both changed the same row, then t2 wins the conflict and t1 aborts. In other words, t1 was trying to change data that had been rendered stale by t2.

Note

If two transactions are bound to conflict more often than not, then it is a good practice to start them on the same server. They then have a chance to synchronize on the local lock manager instead of aborting later in the replication protocol.

21.9.4 Data Definition Statements

Care needs to be taken when executing data definition statements also commonly known as data definition language (DDL). Given that MySQL does not support atomic or transactional DDL, one cannot optimistically execute DDL statements and later roll back if needs be. Consequently, the lack of atomicity does not fit directly into the optimistic replication paradigm that Group Replication is based on.

Therefore, more care needs to be taken when replicating data definition statements. Schema changes and changes to the data that the object contains need to be handled through the same server while the schema operation has not yet completed and replicated everywhere. Failure to do so, may result in data inconsistency.

Note

If the group is deployed in single-primary mode, then this is not a problem, because all changes are performed through the same server, the primary.

Warning

MySQL DDL execution is not atomic or transactional. The server executes and commits without securing group agreement first. As such, you must route DDL and DML for the same object through the same server, while the DDL is executing and has not replicated everywhere yet.

21.9.5 Distributed Recovery

This section describes the process through which a joining member catches up with the remaining servers, called the distributed recovery phase.

21.9.5.1 Group Replication Basics

The Group Replication distributed recovery procedure can be summarized as the process through which a new server gets missing data from a online server in the group, while listening for events happening in the group. During recovery a server listens to membership events and also to the transactions that are happening while recovery is happening. This is a high level summary. The following sections provide additional detail, by describing the two phases of the procedure.

Phase 1

In the first phase, the joiner (joining server), will select one of the online servers on the group to be the *donor* of the state that it is missing. The donor is responsible for handing the joiner all the data it is missing up to the moment it has joined the group. This is achieved by relying on a standard asynchronous replication channel, established between donor and joiner. Through that channel, binary logs flow up until the point that the view change happened when the joiner became part of the group. While the joiner receives binary logs from the donor, it is also applying them.

Furthermore, while the binary logs transfer is ongoing, the joiner is also caching every transaction that is exchanged within the group. I.e., it is listening for transactions that are happening after it joined the group and while it is applying the missing state from the donor. When the first phase ends and the replication channel to the donor is closed, the joiner then starts phase two: the catch up.

Phase 2

In this phase, the joiner proceeds to the execution of the cached transactions and when the number of transactions queued for execution finally reaches zero, then the member is declared online.

Resilience

The recovery procedure withstands donor failures while the joiner is fetching binary logs from it. In such cases, whenever a donor fails during phase 1, the joiner fails over to a new donor and resumes from that one. When that happens the joiner closes the connection to the failed joiner explicitly and opens a connection to a new donor. This happens automatically.

21.9.5.2 Recovering From a Point-in-time

To synchronize the joiner with the donor up to a specific point in time, the joiner and donor make use of the MySQL global transaction identifiers (GTIDs) mechanism. However, that is not enough, since that only provides a mean to realize which transactions the joiner is missing. It does not help marking a specific point in time to which the joiner must catch up, nor does it help conveying certification information. This is where the binary log view markers come into play. They mark view changes in the binary log stream, and also contain additional metadata information, provisioning the joiner with missing certification related data.

View and View Changes

To explain the concept of view change markers, it is important to understand what a view and a view change are.

A *view* corresponds to a group of members participating actively in the current configuration, in other words at a specific point in time. They are correct and online in the system.

A *view change* occurs when a modification to the group configuration happens, such as a member joining or leaving. Any group membership change results in an independent view change communicated to all members at the same logical point in time.

A *view identifier* uniquely identifies a view. It is generated whenever a view change happens

At the group communication layer, view changes with their associated view ids are then boundaries between the data exchanged before and after a member joins. This concept is implemented through a new binary log event: the"view change log event". The view id thus becomes a marker as well for transactions transmitted before and after changes happen in the group membership.

The view identifier itself is built from two parts: *(i)* one that is randomly generated and *(ii)* a monotonically increasing integer. The first part is generated when the group is created, and remains unchanged while there is at least one member in the group. The second part is incremented every time a view change happens.

The reason for this heterogeneous pair that makes up the view id is the need to unambiguously mark group changes whenever a member joins or leaves but also whenever all members leave the group and no information remains of what view the group was in. In fact, the sole use of monotonic increasing identifiers could lead to the reuse of the same id after full group shutdowns, destroying the uniqueness of the binary log data markers that recovery depends on. To summarize, the first part identifies whenever the group was started from the beginning and the incremental part when the group changed from that point on.

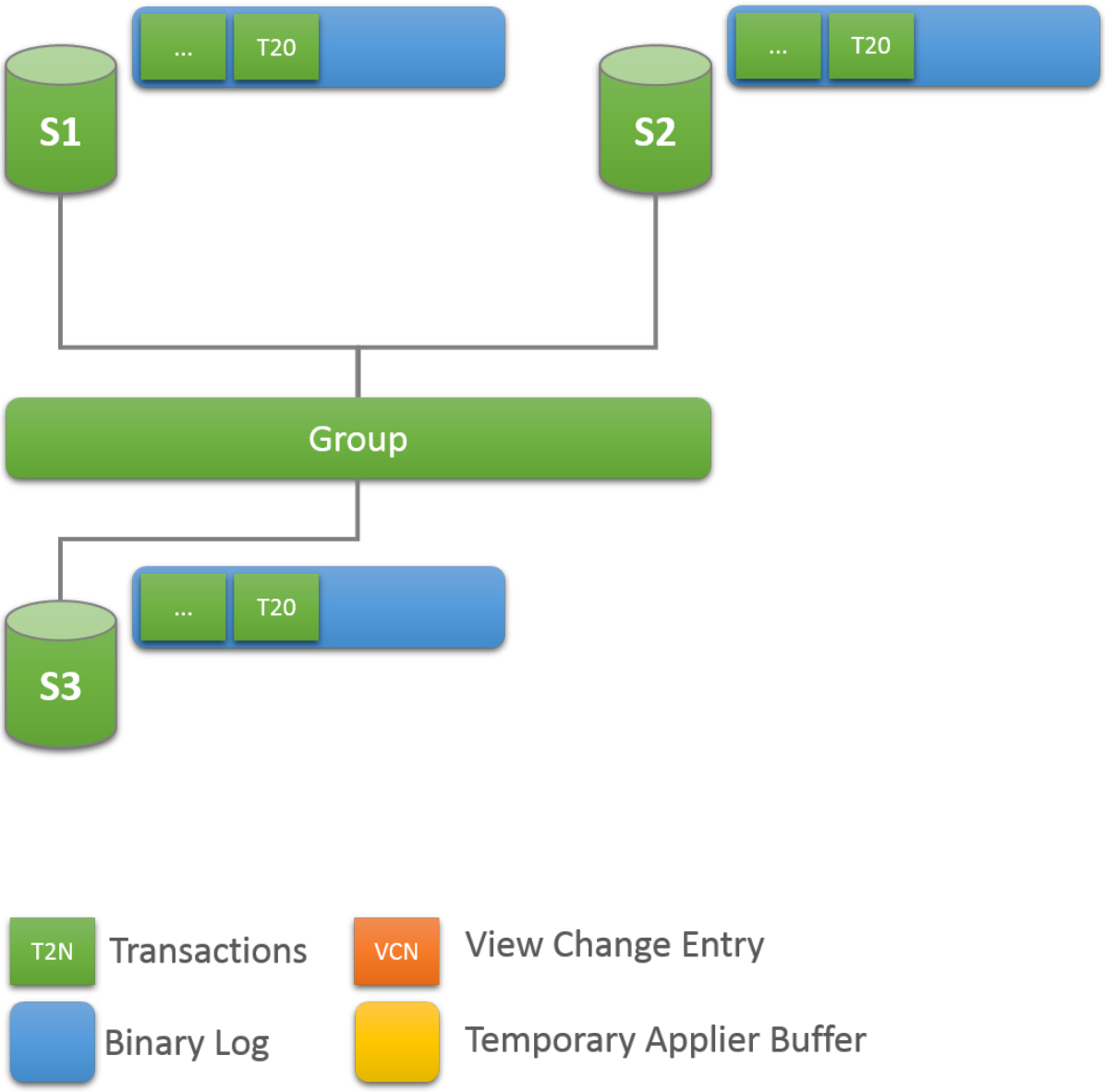
21.9.5.3 View Changes

This section explains the process which controls how the view change identifier is incorporated into a binary log event and written to the log, The following steps are taken:

Begin: Stable Group

All servers are online and processing incoming transactions from the group. Some servers may be a little behind in terms of transactions replicated, but eventually they converge. The group acts as one distributed and replicated database.

Figure 21.10 Stable Group

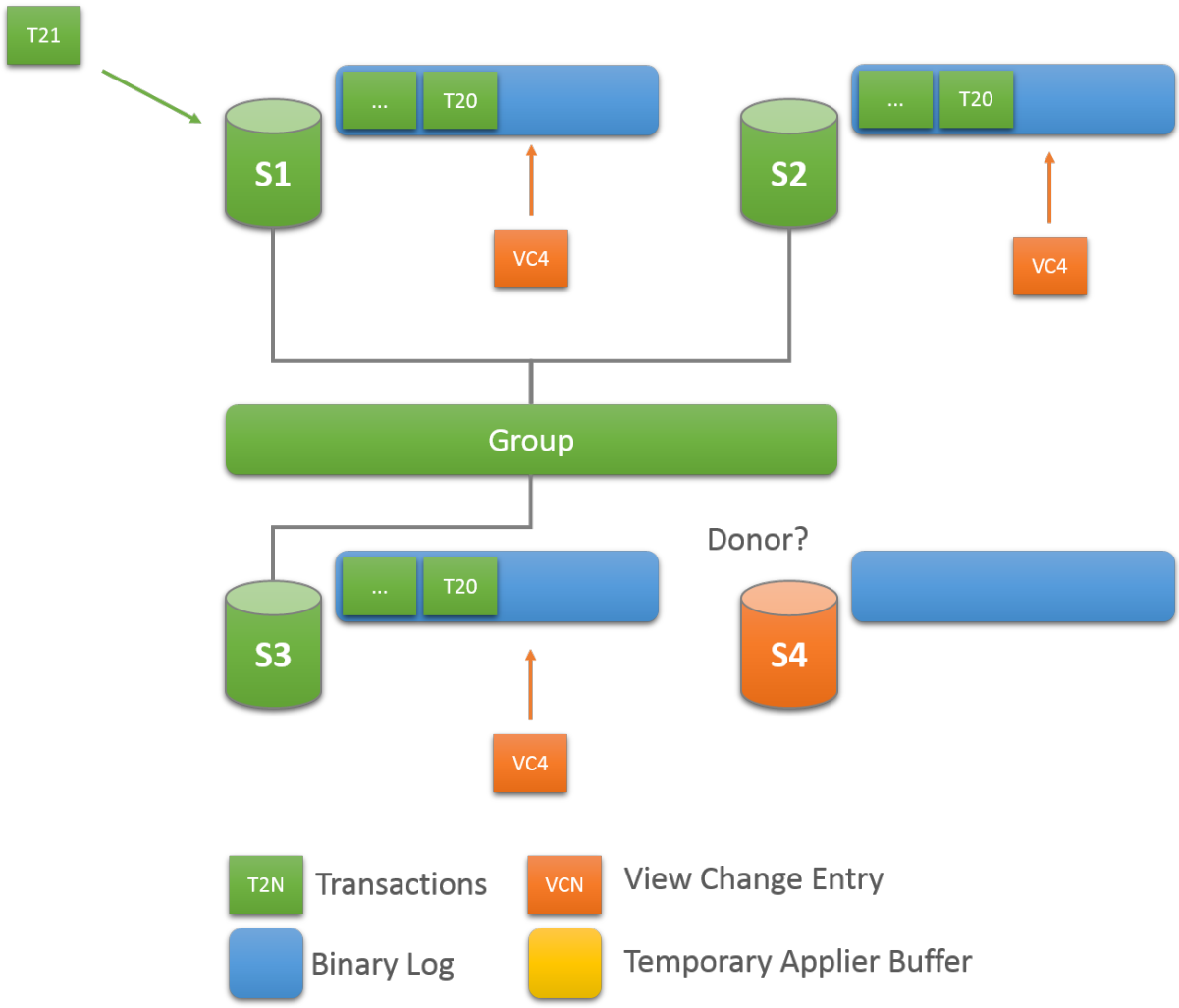


View Change: a Member Joins

Whenever a new member joins the group and therefore a view change is performed, every online server queues a view change log event for execution. This is queued because before the view change, several transactions can be queued on the server to be applied and as such, these belong to the old view. Queuing the view change event after them guarantees a correct marking of when this happened.

Meanwhile, the joiner selects the donor from the list of online servers as stated by the membership service through the view abstraction. A member joins on view 4 and the online members write a View change event to the binary log.

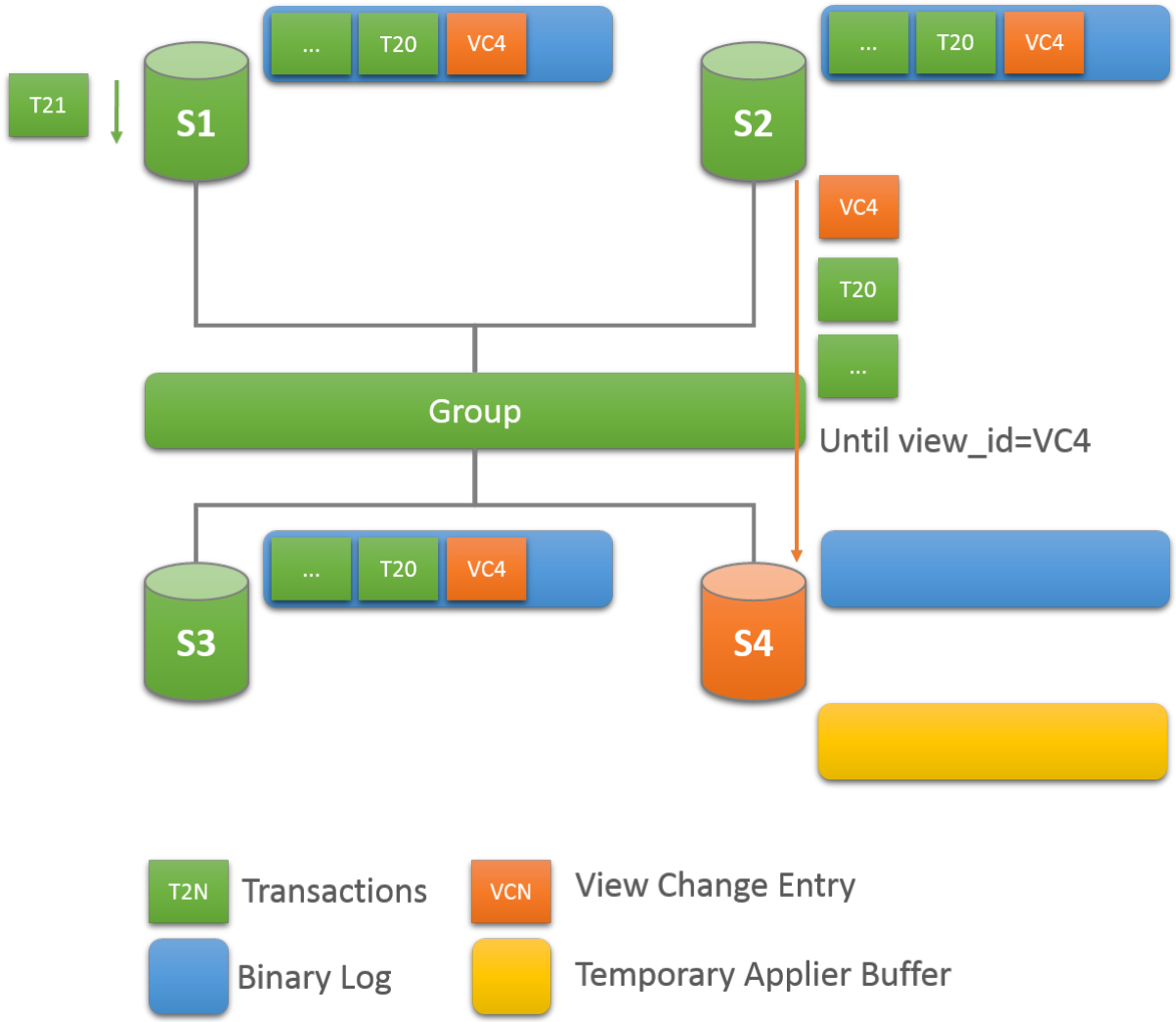
Figure 21.11 A Member Joins



State Transfer: Catching Up

Once the joiner has chosen which server in the group is to be the donor, a new asynchronous replication connection is established between the two and the state transfer begins (phase 1). This interaction with the donor continues until the joiner's applier thread processes the view change log event that corresponds to the view change triggered when the joiner came into the group. In other words, the joiner replicates from the donor, until it gets to the marker with the view identifier which matches the view marker it is already in.

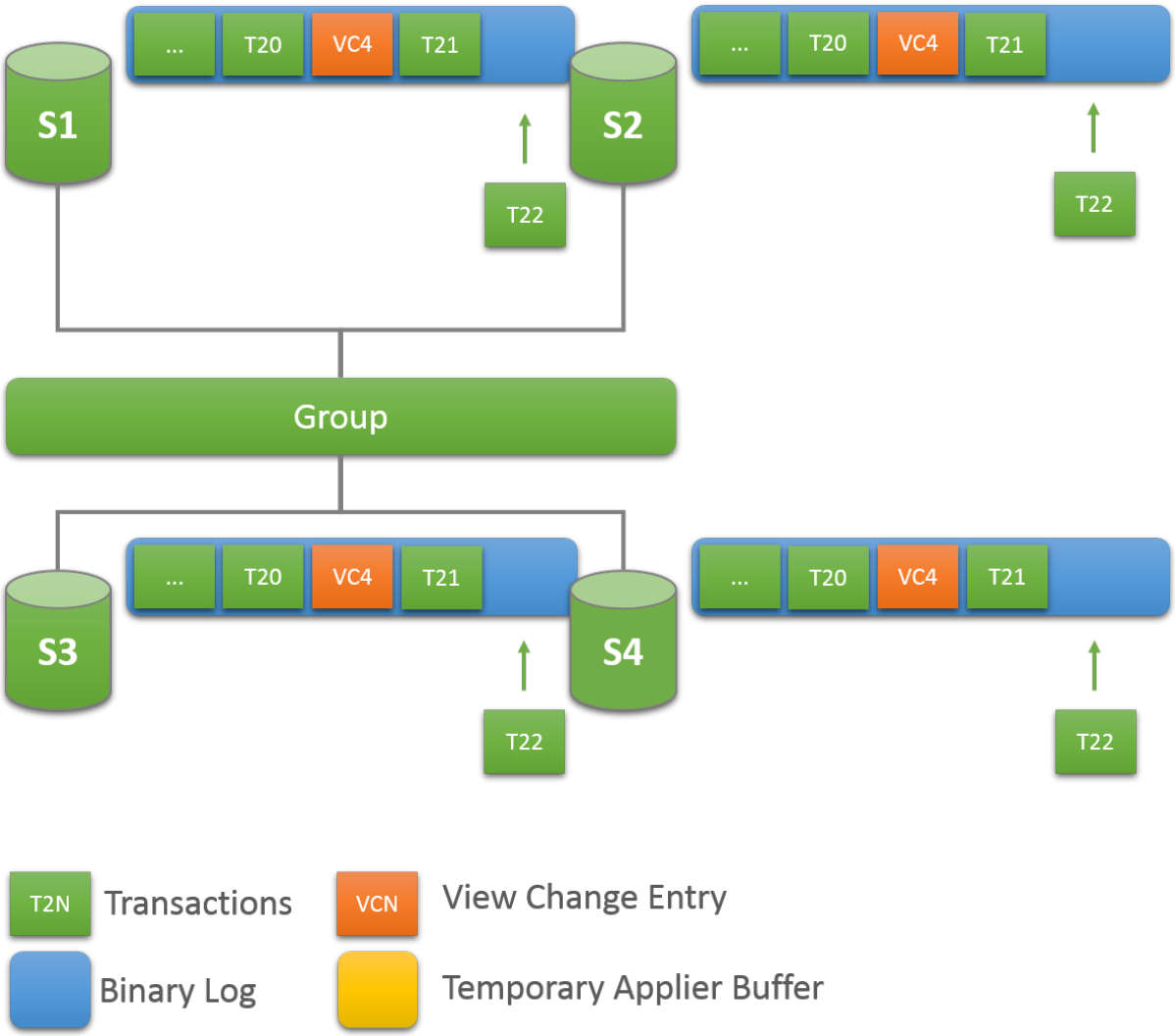
Figure 21.12 State Transfer: Catching Up



As view identifiers are transmitted to all members in the group at the same logical time, the joiner knows at which view identifier it should stop replicating. This avoids complex GTID set calculations because the view id clearly marks which data belongs to each group view.

While the joiner is replicating from the donor, it is also caching incoming transactions from the group. Eventually, it stops replicating from the donor and switches to applying those that are cached.

Figure 21.13 Queued Transactions



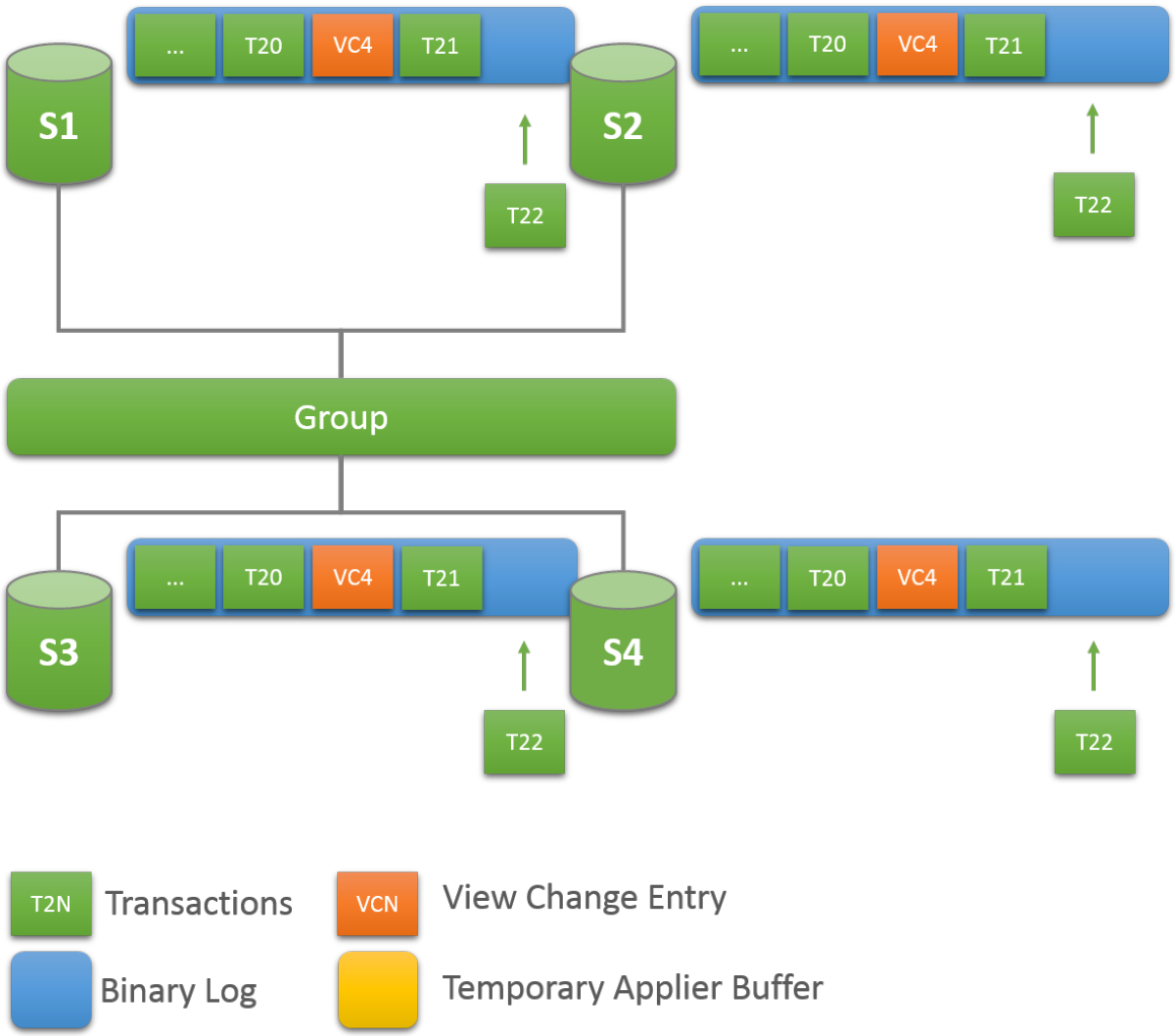
Finish: Caught Up

When the joiner recognizes a view change log event with the expected view identifier, the connection to the donor is terminated and it starts applying the cached transactions. An important point to understand is the final recovery procedure. Although it acts as a marker in the binary log, delimiting view changes, the view change log event also plays another role. It conveys the certification information as perceived by all servers when the joiner entered the group, in other words the last view change. Without it, the joiner would not have the necessary information to be able to certify (detect conflicts) subsequent transactions.

The duration of the catch up (phase 2) is not deterministic, because it depends on the workload and the rate of incoming transactions to the group. This process is completely online and the joiner does not block any other server in the group while it is catching up. Therefore the number of transactions the joiner is behind when it moves to phase 2 can, for this reason, vary and thus increase or decrease according to the workload.

When the joiner's queued transactions reach zero and its stored data is equal to the other members, its public state changes to online.

Figure 21.14 Instance Online



21.9.5.4 Usage Advice and Limitations of Distributed Recovery

Distributed recovery does have some limitations. It is based on classic asynchronous replication and as such it may be slow if the joiner is not provisioned at all or is provisioned with a very old backup image. This means that if the data to transfer is too big at phase 1, the server may take a very long time to recover. As such, the recommendation is that before adding a server to the group, one should provision it with a fairly recent snapshot of a server already in the group. This minimizes the length of phase 1 and reduces the impact on the donor server, since it has to save and transfer less binary logs.

Warning

It is recommended that a server is provisioned before it is added to a group. That way, one minimizes the time spent on the recovery step.

21.9.6 Observability

There is a lot of automation built into the Group Replication plugin. Nonetheless, you might sometimes need to understand what is happening behind the scenes. This is where the instrumentation of Group Replication and Performance Schema becomes important. The entire state of the system (including the view, conflict statistics and service states) can be queried through performance_schema tables. The distributed nature of the replication protocol and the fact that server instances agree and thus synchronize on transactions and metadata makes it simpler to inspect the state of the group. For

example, you can connect to a single server in the group and obtain both local and global information by issuing select statements on the Group Replication related Performance Schema tables. For more information, see Section 21.3, “Monitoring Group Replication”.

21.9.7 Group Replication Performance

This section explains how to use the available configuration options to gain the best performance from your group.

21.9.7.1 Fine Tuning the Group Communication Thread

The group communication thread (GCT) runs in a loop while the Group Replication plugin is loaded. The GCT receives messages from the group and from the plugin, handles quorum and failure detection related tasks, sends out some keep alive messages and also handles the incoming and outgoing transactions from/to the server/group. The GCT waits for incoming messages in a queue. When there are no messages, the GCT waits. By configuring this wait to be a little longer (doing an active wait) before actually going to sleep can prove to be beneficial in some cases. This is because the alternative is for the operating system to switch out the GCT from the processor and do a context switch.

To force the GCT do an active wait, use the `group_replication_poll_spin_loops` option, which makes the GCT loop, doing nothing relevant for the configured number of loops, before actually polling the queue for the next message.

For example:

```
mysql> SET GLOBAL group_replication_poll_spin_loops= 10000;
```

21.9.7.2 Message Compression

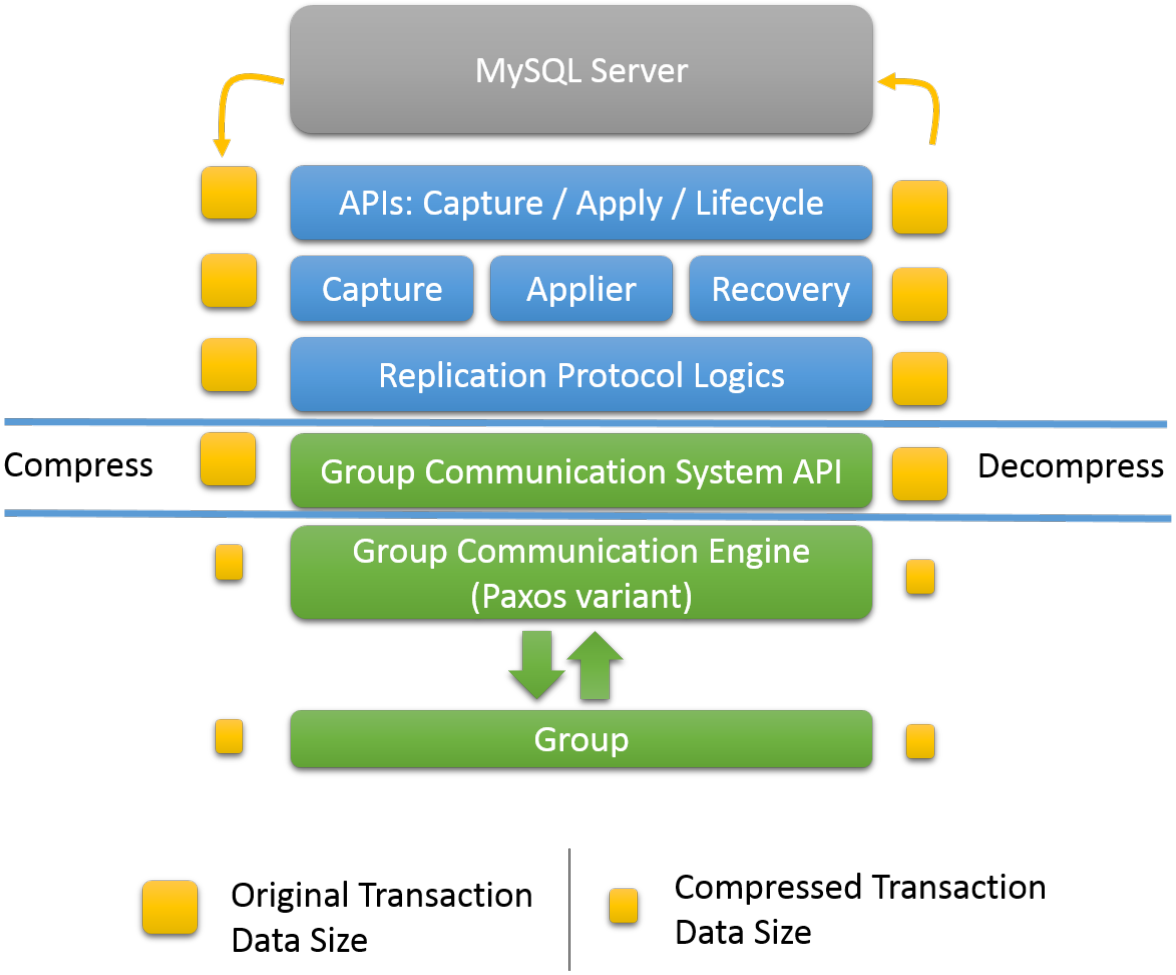
When network bandwidth is a bottleneck, message compression can provide up to 30-40% throughput improvement at the group communication level. This is especially important within the context of large groups of servers under load.

Table 21.8 LZ4 Compression Ratios for Different Binary Log Formats.

	Ratio	
Workload	**ROW**	**STMT**
mysqlslapd	4,5	4,1
sysbench	3,4	2,9

The TCP peer-to-peer nature of the interconnections between *N* participants on the group makes the sender send the same amount of data *N* times. Furthermore, binary logs are likely to exhibit a high compression ratio (see table above). This makes compression a compelling feature for workloads that contain large transaction.

Figure 21.15 Compression Support



Compression happens at the group communication engine level, before the data is handed over to the group communication thread, so it happens within the context of the mysql user session thread. Transaction payloads may be compressed before being sent out to the group and decompressed when received. Compression is conditional and depends on a configured threshold. By default compression is enabled.

In addition, there is no requirement that all servers in the group have compression enabled to be able to work together. Upon receiving a message, the member checks the message envelope to verify whether it is compressed or not. If needed, then the member decompresses the transaction, before delivering it to the upper layer.

The compression algorithm used is LZ4. Compression is enabled by default with threshold of 1000000 bytes. The compression threshold, in bytes, may be set to something larger than default. In that case, only transactions that have a payload larger than the threshold are compressed. Below is an example of how to set a compression threshold.

```
STOP GROUP_REPLICATION;
SET GLOBAL group_replication_compression_threshold= 2097152;
START GROUP_REPLICATION;
```

This sets the compression threshold to 2MB. If a transaction generates a replication message with a payload larger than 2MB, for example a binary log transaction entry larger than 2MB, then it is compressed. To disable compression set threshold to 0.

21.9.7.3 Flow Control

Group Replication ensures that a transaction only commits after a majority of the members in a group have received it and agreed on the relative order between all transactions that were sent concurrently.

This approach works well if the total number of writes to the group does not exceed the write capacity of any member in the group. If it does and some of the members have less write throughput than others, particularly less than the writer members, those members can start lagging behind of the writers.

Having some members lagging behind the group brings some problematic consequences, particularly, the reads on such members may externalize very old data. Depending on why the member is lagging behind, other members in the group may have to save more or less replication context to be able to fulfil potential data transfer requests from the slow member.

There is however a mechanism in the replication protocol to avoid having too much distance, in terms of transactions applied, between fast and slow members. This is known as the flow control mechanism. It tries to address several goals:

1. to keep the members close enough to make buffering and de-synchronization between members a small problem;
2. to adapt quickly to changing conditions like different workloads or more writers in the group;
3. to give each member a fair share of the available write capacity;
4. to not reduce throughput more than strictly necessary to avoid wasting resources.

Given the design of Group Replication, the decision whether to throttle or not may be decided taking into account two work queues: *(i)* the *certification* queue; *(ii)* and on the binary log *applier* queue. Whenever the size of one of these queues exceeds the user-defined threshold, the throttling mechanism is triggered. Only configure: *(i)* whether to do flow control at the certifier or at the applier level, or both; and *(ii)* what is the threshold for each queue.

The flow control depends on two basic mechanisms:

1. the monitoring of members to collect some statistics on throughput and queue sizes of all group members to make educated guesses on what is the maximum write pressure each member should be subjected to;
2. the throttling of members that are trying to write beyond their fair-share of the available capacity at each moment in time.

Probes and Statistics

The monitoring mechanism works by having each member deploying a set of probes to collect information about its work queues and throughput. It then propagates that information to the group periodically to share that data with the other members.

Such probes are scattered throughout the plugin stack and allow one to establish metrics, such as:

- the certifier queue size;
- the replication applier queue size;
- the total number of transactions certified;
- the total number of remote transactions applied in the member;
- the total number of local transactions.

Once a member receives a message with statistics from another member, it calculates additional metrics regarding how many transactions were certified, applied and locally executed in the last monitoring period.

Monitoring data is shared with others in the group periodically. The monitoring period must be high enough to allow the other members to decide on the current write requests, but low enough that it has minimal impact on group bandwidth. The information is shared every second, and this period is sufficient to address both concerns.

Group Replication Throttling

Based on the metrics gathered across all servers in the group, a throttling mechanism kicks in and decides whether to limit the rate a member is able to execute/commit new transactions.

Therefore, metrics acquired from all members are the basis for calculating the capacity of each member: if a member has a large queue (for certification or the applier thread), then the capacity to execute new transactions should be close to ones certified or applied in the last period.

The lowest capacity of all the members in the group determines the real capacity of the group, while the number of local transactions determines how many members are writing to it, and, consequently, how many members should that available capacity be shared with.

This means that every member has an established write quota based on the available capacity, in other words a number of transactions it can safely issue for the next period. The writer-quota will be enforced by the throttling mechanism if the queue size of the certifier or the binary log applier exceeds a user-defined threshold.

The quota is reduced by the number of transactions that were delayed in the last period, and then also further reduced by 10% to allow the queue that triggered the problem to reduce its size. In order to avoid large jumps in throughput once the queue size goes beyond the threshold, the throughput is only allowed to grow by the same 10% per period after that.

The current throttling mechanism does not penalize transactions below quota, but delays finishing those transactions that exceed it until the end of the monitoring period. As a consequence, if the quota is very small for the write requests issued some transactions may have latencies close to the monitoring period.

Chapter 22 Partitioning

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This chapter discusses *user-defined partitioning*. In MySQL 8.0, partitioning support is provided by the `InnoDB` storage engine. (The `NDB` storage engine used by MySQL Cluster also provides partitioning support, but `NDB` is not included in MySQL 8.0.)

MySQL 8.0 does not currently support partitioning of tables using any storage engine other than `InnoDB`, such as `MyISAM`. An attempt to create a partitioned tables using a storage engine that does not supply native partitioning support fails with `ER_CHECK_NOT_IMPLEMENTED`.

MySQL 8.0 Community binaries provided by Oracle include partitioning support provided by the `InnoDB` storage engine. For information about partitioning support offered in MySQL Enterprise Edition binaries, see Chapter 29, *MySQL Enterprise Edition*.

If you are compiling MySQL 8.0 from source, configuring the build with `InnoDB` support is sufficient to produce binaries with partition support for `InnoDB` tables. For more information, see Section 2.8, “Installing MySQL from Source”.

Nothing further needs to be done to enable partitioning support by `InnoDB` (for example, no special entries are required in the `my.cnf` file).

It is not possible to disable partitioning support by the `InnoDB` storage engine.

See Section 22.1, “Overview of Partitioning in MySQL”, for an introduction to partitioning and partitioning concepts.

Several types of partitioning are supported, as well as subpartitioning; see Section 22.2, “Partitioning Types”, and Section 22.2.6, “Subpartitioning”.

Section 22.3, “Partition Management”, covers methods of adding, removing, and altering partitions in existing partitioned tables.

Section 22.3.4, “Maintenance of Partitions”, discusses table maintenance commands for use with partitioned tables.

The `PARTITIONS` table in the `INFORMATION_SCHEMA` database provides information about partitions and partitioned tables. See Section 24.13, “The INFORMATION_SCHEMA PARTITIONS Table”, for more information; for some examples of queries against this table, see Section 22.2.7, “How MySQL Partitioning Handles NULL”.

For known issues with partitioning in MySQL 8.0, see Section 22.6, “Restrictions and Limitations on Partitioning”.

You may also find the following resources to be useful when working with partitioned tables.

Additional Resources. Other sources of information about user-defined partitioning in MySQL include the following:

- MySQL Partitioning Forum

  This is the official discussion forum for those interested in or experimenting with MySQL Partitioning technology. It features announcements and updates from MySQL developers and others. It is monitored by members of the Partitioning Development and Documentation Teams.

- Mikael Ronström's Blog

  MySQL Partitioning Architect and Lead Developer Mikael Ronström frequently posts articles here concerning his work with MySQL Partitioning and MySQL Cluster.

- PlanetMySQL

  A MySQL news site featuring MySQL-related blogs, which should be of interest to anyone using my MySQL. We encourage you to check here for links to blogs kept by those working with MySQL Partitioning, or to have your own blog added to those covered.

22.1 Overview of Partitioning in MySQL

This section provides a conceptual overview of partitioning in MySQL 8.0.

For information on partitioning restrictions and feature limitations, see Section 22.6, “Restrictions and Limitations on Partitioning”.

The SQL standard does not provide much in the way of guidance regarding the physical aspects of data storage. The SQL language itself is intended to work independently of any data structures or media underlying the schemas, tables, rows, or columns with which it works. Nonetheless, most advanced database management systems have evolved some means of determining the physical location to be used for storing specific pieces of data in terms of the file system, hardware or even both. In MySQL, the `InnoDB` storage engine has long supported the notion of a tablespace (see Section 15.7, “InnoDB Tablespaces”), and the MySQL Server, even prior to the introduction of partitioning, could be configured to employ different physical directories for storing different databases (see Section 8.12.2, “Using Symbolic Links”, for an explanation of how this is done).

Partitioning takes this notion a step further, by enabling you to distribute portions of individual tables across a file system according to rules which you can set largely as needed. In effect, different portions of a table are stored as separate tables in different locations. The user-selected rule by which the division of data is accomplished is known as a *partitioning function*, which in MySQL can be the modulus, simple matching against a set of ranges or value lists, an internal hashing function, or a linear hashing function. The function is selected according to the partitioning type specified by the user, and takes as its parameter the value of a user-supplied expression. This expression can be a column value, a function acting on one or more column values, or a set of one or more column values, depending on the type of partitioning that is used.

In the case of `RANGE`, `LIST`, and [`LINEAR`] `HASH` partitioning, the value of the partitioning column is passed to the partitioning function, which returns an integer value representing the number of the partition in which that particular record should be stored. This function must be nonconstant and nonrandom. It may not contain any queries, but may use an SQL expression that is valid in MySQL, as long as that expression returns either `NULL` or an integer *`intval`* such that

```
-MAXVALUE <= intval <= MAXVALUE
```

(`MAXVALUE` is used to represent the least upper bound for the type of integer in question. `-MAXVALUE` represents the greatest lower bound.)

For [`LINEAR`] `KEY`, `RANGE COLUMNS`, and `LIST COLUMNS` partitioning, the partitioning expression consists of a list of one or more columns.

For [`LINEAR`] `KEY` partitioning, the partitioning function is supplied by MySQL.

For more information about permitted partitioning column types and partitioning functions, see Section 22.2, “Partitioning Types”, as well as Section 13.1.15, “CREATE TABLE Syntax”, which provides partitioning syntax descriptions and additional examples. For information about restrictions on partitioning functions, see Section 22.6.3, “Partitioning Limitations Relating to Functions”.

This is known as *horizontal partitioning*—that is, different rows of a table may be assigned to different physical partitions. MySQL 8.0 does not support *vertical partitioning*, in which different columns of a table are assigned to different physical partitions. There are no plans at this time to introduce vertical partitioning into MySQL.

For creating partitioned tables, you must use a storage engine that supports them. In MySQL 8.0, all partitions of the same partitioned table must use the same storage engine. However, there is nothing preventing you from using different storage engines for different partitioned tables on the same MySQL server or even in the same database.

In MySQL 8.0, the only storage engine that supports partitioning is `InnoDB`. Partitioning cannot be used with storage engines that do not support it; these include the `MyISAM`, `MERGE`, `CSV`, and `FEDERATED` storage engines.

When creating a partitioned table, the default storage engine is used just as when creating any other table; to override this behavior, it is necessary only to use the `[STORAGE] ENGINE` option just as you would for a table that is not partitioned. The target storage engine must provide native partitioning support, or the statement fails. You should keep in mind that `[STORAGE] ENGINE` (and other table options) need to be listed *before* any partitioning options are used in a `CREATE TABLE` statement. This example shows how to create a table that is partitioned by hash into 6 partitions and which uses the `InnoDB` storage engine (regardless of the value of `default_storage_engine`):

```
CREATE TABLE ti (id INT, amount DECIMAL(7,2), tr_date DATE)
    ENGINE=INNODB
    PARTITION BY HASH( MONTH(tr_date) )
    PARTITIONS 6;
```

Each `PARTITION` clause can include a `[STORAGE] ENGINE` option, but in MySQL 8.0 this has no effect.

Unless otherwise specified, the remaining examples in this discussion assume that `default_storage_engine` is `InnoDB`.

Important

Partitioning applies to all data and indexes of a table; you cannot partition only the data and not the indexes, or vice versa, nor can you partition only a portion of the table.

Data and indexes for each partition can be assigned to a specific directory using the `DATA DIRECTORY` and `INDEX DIRECTORY` options for the `PARTITION` clause of the `CREATE TABLE` statement used to create the partitioned table.

InnoDB supports the `DATA DIRECTORY` and `INDEX DIRECTORY` options for individual partitions and subpartitions.

All columns used in the table's partitioning expression must be part of every unique key that the table may have, including any primary key. This means that a table such as this one, created by the following SQL statement, cannot be partitioned:

```
CREATE TABLE tnp (
    id INT NOT NULL AUTO_INCREMENT,
    ref BIGINT NOT NULL,
    name VARCHAR(255),
    PRIMARY KEY pk (id),
    UNIQUE KEY uk (name)
);
```

Because the keys `pk` and `uk` have no columns in common, there are no columns available for use in a partitioning expression. Possible workarounds in this situation include adding the `name` column to the table's primary key, adding the `id` column to `uk`, or simply removing the unique key altogether. See Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”, for more information.

In addition, `MAX_ROWS` and `MIN_ROWS` can be used to determine the maximum and minimum numbers of rows, respectively, that can be stored in each partition. See Section 22.3, “Partition Management”, for more information on these options.

Some advantages of partitioning are listed here:

- Partitioning makes it possible to store more data in one table than can be held on a single disk or file system partition.
- Data that loses its usefulness can often be easily removed from a partitioned table by dropping the partition (or partitions) containing only that data. Conversely, the process of adding new data can in some cases be greatly facilitated by adding one or more new partitions for storing specifically that data.
- Some queries can be greatly optimized in virtue of the fact that data satisfying a given `WHERE` clause can be stored only on one or more partitions, which automatically excludes any remaining partitions from the search. Because partitions can be altered after a partitioned table has been created, you can reorganize your data to enhance frequent queries that may not have been often used when the partitioning scheme was first set up. This ability to exclude non-matching partitions (and thus any rows they contain) is often referred to as *partition pruning*. For more information, see Section 22.4, “Partition Pruning”.

  In addition, MySQL 8.0 supports explicit partition selection for queries. For example, `SELECT * FROM t PARTITION (p0,p1) WHERE c < 5` selects only those rows in partitions `p0` and `p1` that match the `WHERE` condition. In this case, MySQL does not check any other partitions of table `t`; this can greatly speed up queries when you already know which partition or partitions you wish to examine. Partition selection is also supported for the data modification statements `DELETE`, `INSERT`, `REPLACE`, `UPDATE`, and `LOAD DATA`, `LOAD XML`. See the descriptions of these statements for more information and examples.

Other benefits usually associated with partitioning include those in the following list. These features are not currently implemented in MySQL Partitioning, but are high on our list of priorities.

- Queries involving aggregate functions such as `SUM()` and `COUNT()` can easily be parallelized. A simple example of such a query might be `SELECT salesperson_id, COUNT(orders) as order_total FROM sales GROUP BY salesperson_id;`. By “parallelized,” we mean that the

query can be run simultaneously on each partition, and the final result obtained merely by summing the results obtained for all partitions.

- Achieving greater query throughput in virtue of spreading data seeks over multiple disks.

22.2 Partitioning Types

This section discusses the types of partitioning which are available in MySQL 8.0. These include the types listed here:

- **RANGE partitioning.** This type of partitioning assigns rows to partitions based on column values falling within a given range. See Section 22.2.1, “RANGE Partitioning”. For information about an extension to this type, `RANGE COLUMNS`, see Section 22.2.3.1, “RANGE COLUMNS partitioning”.
- **LIST partitioning.** Similar to partitioning by `RANGE`, except that the partition is selected based on columns matching one of a set of discrete values. See Section 22.2.2, “LIST Partitioning”. For information about an extension to this type, `LIST COLUMNS`, see Section 22.2.3.2, “LIST COLUMNS partitioning”.
- **HASH partitioning.** With this type of partitioning, a partition is selected based on the value returned by a user-defined expression that operates on column values in rows to be inserted into the table. The function may consist of any expression valid in MySQL that yields a nonnegative integer value. An extension to this type, `LINEAR HASH`, is also available. See Section 22.2.4, “HASH Partitioning”.
- **KEY partitioning.** This type of partitioning is similar to partitioning by `HASH`, except that only one or more columns to be evaluated are supplied, and the MySQL server provides its own hashing function. These columns can contain other than integer values, since the hashing function supplied by MySQL guarantees an integer result regardless of the column data type. An extension to this type, `LINEAR KEY`, is also available. See Section 22.2.5, “KEY Partitioning”.

A very common use of database partitioning is to segregate data by date. Some database systems support explicit date partitioning, which MySQL does not implement in 8.0. However, it is not difficult in MySQL to create partitioning schemes based on `DATE`, `TIME`, or `DATETIME` columns, or based on expressions making use of such columns.

When partitioning by `KEY` or `LINEAR KEY`, you can use a `DATE`, `TIME`, or `DATETIME` column as the partitioning column without performing any modification of the column value. For example, this table creation statement is perfectly valid in MySQL:

```
CREATE TABLE members (
    firstname VARCHAR(25) NOT NULL,
    lastname VARCHAR(25) NOT NULL,
    username VARCHAR(16) NOT NULL,
    email VARCHAR(35),
    joined DATE NOT NULL
)
PARTITION BY KEY(joined)
PARTITIONS 6;
```

In MySQL 8.0, it is also possible to use a `DATE` or `DATETIME` column as the partitioning column using `RANGE COLUMNS` and `LIST COLUMNS` partitioning.

Other partitioning types require a partitioning expression that yields an integer value or `NULL`. If you wish to use date-based partitioning by `RANGE`, `LIST`, `HASH`, or `LINEAR HASH`, you can simply employ a function that operates on a `DATE`, `TIME`, or `DATETIME` column and returns such a value, as shown here:

```
CREATE TABLE members (
    firstname VARCHAR(25) NOT NULL,
    lastname VARCHAR(25) NOT NULL,
    username VARCHAR(16) NOT NULL,
    email VARCHAR(35),
    joined DATE NOT NULL
)
PARTITION BY RANGE( YEAR(joined) ) (
    PARTITION p0 VALUES LESS THAN (1960),
    PARTITION p1 VALUES LESS THAN (1970),
    PARTITION p2 VALUES LESS THAN (1980),
    PARTITION p3 VALUES LESS THAN (1990),
    PARTITION p4 VALUES LESS THAN MAXVALUE
);
```

Additional examples of partitioning using dates may be found in the following sections of this chapter:

- Section 22.2.1, “RANGE Partitioning”
- Section 22.2.4, “HASH Partitioning”
- Section 22.2.4.1, “LINEAR HASH Partitioning”

For more complex examples of date-based partitioning, see the following sections:

- Section 22.4, “Partition Pruning”
- Section 22.2.6, “Subpartitioning”

MySQL partitioning is optimized for use with the `TO_DAYS()`, `YEAR()`, and `TO_SECONDS()` functions. However, you can use other date and time functions that return an integer or `NULL`, such as `WEEKDAY()`, `DAYOFYEAR()`, or `MONTH()`. See Section 12.7, “Date and Time Functions”, for more information about such functions.

It is important to remember—regardless of the type of partitioning that you use—that partitions are always numbered automatically and in sequence when created, starting with `0`. When a new row is inserted into a partitioned table, it is these partition numbers that are used in identifying the correct partition. For example, if your table uses 4 partitions, these partitions are numbered `0`, `1`, `2`, and `3`. For the `RANGE` and `LIST` partitioning types, it is necessary to ensure that there is a partition defined for each partition number. For `HASH` partitioning, the user-supplied expression must evaluate to an integer value greater than `0`. For `KEY` partitioning, this issue is taken care of automatically by the hashing function which the MySQL server employs internally.

Names of partitions generally follow the rules governing other MySQL identifiers, such as those for tables and databases. However, you should note that partition names are not case-sensitive. For example, the following `CREATE TABLE` statement fails as shown:

```
mysql> CREATE TABLE t2 (val INT)
    -> PARTITION BY LIST(val)(
    ->     PARTITION mypart VALUES IN (1,3,5),
    ->     PARTITION MyPart VALUES IN (2,4,6)
    -> );
ERROR 1488 (HY000): Duplicate partition name mypart
```

Failure occurs because MySQL sees no difference between the partition names `mypart` and `MyPart`.

When you specify the number of partitions for the table, this must be expressed as a positive, nonzero integer literal with no leading zeros, and may not be an expression such as `0.8E+01` or `6-2`, even if it evaluates to an integer value. Decimal fractions are not permitted.

In the sections that follow, we do not necessarily provide all possible forms for the syntax that can be used for creating each partition type; for this information, see Section 13.1.15, “CREATE TABLE Syntax”.

22.2.1 RANGE Partitioning

A table that is partitioned by range is partitioned in such a way that each partition contains rows for which the partitioning expression value lies within a given range. Ranges should be contiguous but not overlapping, and are defined using the `VALUES LESS THAN` operator. For the next few examples, suppose that you are creating a table such as the following to hold personnel records for a chain of 20 video stores, numbered 1 through 20:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
);
```

Note

The `employees` table used here has no primary or unique keys. While the examples work as shown for purposes of the present discussion, you should keep in mind that tables are extremely likely in practice to have primary keys, unique keys, or both, and that allowable choices for partitioning columns depend on the columns used for these keys, if any are present. For a discussion of these issues, see Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”.

This table can be partitioned by range in a number of ways, depending on your needs. One way would be to use the `store_id` column. For instance, you might decide to partition the table 4 ways by adding a `PARTITION BY RANGE` clause as shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
)
PARTITION BY RANGE (store_id) (
    PARTITION p0 VALUES LESS THAN (6),
    PARTITION p1 VALUES LESS THAN (11),
    PARTITION p2 VALUES LESS THAN (16),
    PARTITION p3 VALUES LESS THAN (21)
);
```

In this partitioning scheme, all rows corresponding to employees working at stores 1 through 5 are stored in partition `p0`, to those employed at stores 6 through 10 are stored in partition `p1`, and so on. Each partition is defined in order, from lowest to highest. This is a requirement of the `PARTITION BY RANGE` syntax; you can think of it as being analogous to a series of `if ... elseif ...` statements in C or Java in this regard.

It is easy to determine that a new row containing the data `(72, 'Mitchell', 'Wilson', '1998-06-25', NULL, 13)` is inserted into partition `p2`, but what happens when your chain adds a 21st store? Under this scheme, there is no rule that covers a row whose `store_id` is greater than 20, so an error results because the server does not know where to place it. You can keep this from occurring by using a “catchall” `VALUES LESS THAN` clause in the `CREATE TABLE` statement that provides for all values greater than the highest value explicitly named:

```
CREATE TABLE employees (
```

```
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
)
PARTITION BY RANGE (store_id) (
    PARTITION p0 VALUES LESS THAN (6),
    PARTITION p1 VALUES LESS THAN (11),
    PARTITION p2 VALUES LESS THAN (16),
    PARTITION p3 VALUES LESS THAN MAXVALUE
);
```

(As with the other examples in this chapter, we assume that the default storage engine is `InnoDB`.)

Note

Another way to avoid an error when no matching value is found is to use the `IGNORE` keyword as part of the `INSERT` statement. For an example, see Section 22.2.2, “LIST Partitioning”. Also see Section 13.2.6, “INSERT Syntax”, for general information about `IGNORE`.

`MAXVALUE` represents an integer value that is always greater than the largest possible integer value (in mathematical language, it serves as a *least upper bound*). Now, any rows whose `store_id` column value is greater than or equal to 16 (the highest value defined) are stored in partition `p3`. At some point in the future—when the number of stores has increased to 25, 30, or more—you can use an `ALTER TABLE` statement to add new partitions for stores 21-25, 26-30, and so on (see Section 22.3, “Partition Management”, for details of how to do this).

In much the same fashion, you could partition the table based on employee job codes—that is, based on ranges of `job_code` column values. For example—assuming that two-digit job codes are used for regular (in-store) workers, three-digit codes are used for office and support personnel, and four-digit codes are used for management positions—you could create the partitioned table using the following statement:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
)
PARTITION BY RANGE (job_code) (
    PARTITION p0 VALUES LESS THAN (100),
    PARTITION p1 VALUES LESS THAN (1000),
    PARTITION p2 VALUES LESS THAN (10000)
);
```

In this instance, all rows relating to in-store workers would be stored in partition `p0`, those relating to office and support staff in `p1`, and those relating to managers in partition `p2`.

It is also possible to use an expression in `VALUES LESS THAN` clauses. However, MySQL must be able to evaluate the expression's return value as part of a `LESS THAN` (`<`) comparison.

Rather than splitting up the table data according to store number, you can use an expression based on one of the two `DATE` columns instead. For example, let us suppose that you wish to partition based on the year that each employee left the company; that is, the value of `YEAR(separated)`. An example of a `CREATE TABLE` statement that implements such a partitioning scheme is shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
)
PARTITION BY RANGE ( YEAR(separated) ) (
    PARTITION p0 VALUES LESS THAN (1991),
    PARTITION p1 VALUES LESS THAN (1996),
    PARTITION p2 VALUES LESS THAN (2001),
    PARTITION p3 VALUES LESS THAN MAXVALUE
);
```

In this scheme, for all employees who left before 1991, the rows are stored in partition `p0`; for those who left in the years 1991 through 1995, in `p1`; for those who left in the years 1996 through 2000, in `p2`; and for any workers who left after the year 2000, in `p3`.

It is also possible to partition a table by `RANGE`, based on the value of a `TIMESTAMP` column, using the `UNIX_TIMESTAMP()` function, as shown in this example:

```
CREATE TABLE quarterly_report_status (
    report_id INT NOT NULL,
    report_status VARCHAR(20) NOT NULL,
    report_updated TIMESTAMP NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP
)
PARTITION BY RANGE ( UNIX_TIMESTAMP(report_updated) ) (
    PARTITION p0 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-01-01 00:00:00') ),
    PARTITION p1 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-04-01 00:00:00') ),
    PARTITION p2 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-07-01 00:00:00') ),
    PARTITION p3 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-10-01 00:00:00') ),
    PARTITION p4 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-01-01 00:00:00') ),
    PARTITION p5 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-04-01 00:00:00') ),
    PARTITION p6 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-07-01 00:00:00') ),
    PARTITION p7 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-10-01 00:00:00') ),
    PARTITION p8 VALUES LESS THAN ( UNIX_TIMESTAMP('2010-01-01 00:00:00') ),
    PARTITION p9 VALUES LESS THAN (MAXVALUE)
);
```

Any other expressions involving `TIMESTAMP` values are not permitted. (See Bug #42849.)

Range partitioning is particularly useful when one or more of the following conditions is true:

- You want or need to delete “old” data. If you are using the partitioning scheme shown previously for the `employees` table, you can simply use `ALTER TABLE employees DROP PARTITION p0;` to delete all rows relating to employees who stopped working for the firm prior to 1991. (See Section 13.1.7, “ALTER TABLE Syntax”, and Section 22.3, “Partition Management”, for more information.) For a table with a great many rows, this can be much more efficient than running a `DELETE` query such as `DELETE FROM employees WHERE YEAR(separated) <= 1990;`.

- You want to use a column containing date or time values, or containing values arising from some other series.

- You frequently run queries that depend directly on the column used for partitioning the table. For example, when executing a query such as `EXPLAIN SELECT COUNT(*) FROM employees WHERE separated BETWEEN '2000-01-01' AND '2000-12-31' GROUP BY store_id;`, MySQL can quickly determine that only partition `p2` needs to be scanned because the remaining partitions cannot contain any records satisfying the `WHERE` clause. See Section 22.4, “Partition Pruning”, for more information about how this is accomplished.

A variant on this type of partitioning is `RANGE COLUMNS` partitioning. Partitioning by `RANGE COLUMNS` makes it possible to employ multiple columns for defining partitioning ranges that apply both to placement of rows in partitions and for determining the inclusion or exclusion of specific partitions

when performing partition pruning. See Section 22.2.3.1, “RANGE COLUMNS partitioning”, for more information.

Partitioning schemes based on time intervals. If you wish to implement a partitioning scheme based on ranges or intervals of time in MySQL 8.0, you have two options:

1. Partition the table by `RANGE`, and for the partitioning expression, employ a function operating on a `DATE`, `TIME`, or `DATETIME` column and returning an integer value, as shown here:

```
CREATE TABLE members (
    firstname VARCHAR(25) NOT NULL,
    lastname VARCHAR(25) NOT NULL,
    username VARCHAR(16) NOT NULL,
    email VARCHAR(35),
    joined DATE NOT NULL
)
PARTITION BY RANGE( YEAR(joined) ) (
    PARTITION p0 VALUES LESS THAN (1960),
    PARTITION p1 VALUES LESS THAN (1970),
    PARTITION p2 VALUES LESS THAN (1980),
    PARTITION p3 VALUES LESS THAN (1990),
    PARTITION p4 VALUES LESS THAN MAXVALUE
);
```

   In MySQL 8.0, it is also possible to partition a table by `RANGE` based on the value of a `TIMESTAMP` column, using the `UNIX_TIMESTAMP()` function, as shown in this example:

```
CREATE TABLE quarterly_report_status (
    report_id INT NOT NULL,
    report_status VARCHAR(20) NOT NULL,
    report_updated TIMESTAMP NOT NULL DEFAULT CURRENT_TIMESTAMP ON UPDATE CURRENT_TIMESTAMP
)
PARTITION BY RANGE ( UNIX_TIMESTAMP(report_updated) ) (
    PARTITION p0 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-01-01 00:00:00') ),
    PARTITION p1 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-04-01 00:00:00') ),
    PARTITION p2 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-07-01 00:00:00') ),
    PARTITION p3 VALUES LESS THAN ( UNIX_TIMESTAMP('2008-10-01 00:00:00') ),
    PARTITION p4 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-01-01 00:00:00') ),
    PARTITION p5 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-04-01 00:00:00') ),
    PARTITION p6 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-07-01 00:00:00') ),
    PARTITION p7 VALUES LESS THAN ( UNIX_TIMESTAMP('2009-10-01 00:00:00') ),
    PARTITION p8 VALUES LESS THAN ( UNIX_TIMESTAMP('2010-01-01 00:00:00') ),
    PARTITION p9 VALUES LESS THAN (MAXVALUE)
);
```

   In MySQL 8.0, any other expressions involving `TIMESTAMP` values are not permitted. (See Bug #42849.)

   Note

   It is also possible in MySQL 8.0 to use `UNIX_TIMESTAMP(timestamp_column)` as a partitioning expression for tables that are partitioned by `LIST`. However, it is usually not practical to do so.

2. Partition the table by `RANGE COLUMNS`, using a `DATE` or `DATETIME` column as the partitioning column. For example, the `members` table could be defined using the `joined` column directly, as shown here:

```
CREATE TABLE members (
    firstname VARCHAR(25) NOT NULL,
    lastname VARCHAR(25) NOT NULL,
    username VARCHAR(16) NOT NULL,
    email VARCHAR(35),
    joined DATE NOT NULL
```

```
)
PARTITION BY RANGE COLUMNS(joined) (
    PARTITION p0 VALUES LESS THAN ('1960-01-01'),
    PARTITION p1 VALUES LESS THAN ('1970-01-01'),
    PARTITION p2 VALUES LESS THAN ('1980-01-01'),
    PARTITION p3 VALUES LESS THAN ('1990-01-01'),
    PARTITION p4 VALUES LESS THAN MAXVALUE
);
```

Note

The use of partitioning columns employing date or time types other than `DATE` or `DATETIME` is not supported with `RANGE COLUMNS`.

22.2.2 LIST Partitioning

List partitioning in MySQL is similar to range partitioning in many ways. As in partitioning by `RANGE`, each partition must be explicitly defined. The chief difference between the two types of partitioning is that, in list partitioning, each partition is defined and selected based on the membership of a column value in one of a set of value lists, rather than in one of a set of contiguous ranges of values. This is done by using `PARTITION BY LIST(expr)` where *`expr`* is a column value or an expression based on a column value and returning an integer value, and then defining each partition by means of a `VALUES IN (value_list)`, where *`value_list`* is a comma-separated list of integers.

Note

In MySQL 8.0, it is possible to match against only a list of integers (and possibly `NULL`—see Section 22.2.7, “How MySQL Partitioning Handles NULL”) when partitioning by `LIST`.

However, other column types may be used in value lists when employing `LIST COLUMN` partitioning, which is described later in this section.

Unlike the case with partitions defined by range, list partitions do not need to be declared in any particular order. For more detailed syntactical information, see Section 13.1.15, “CREATE TABLE Syntax”.

For the examples that follow, we assume that the basic definition of the table to be partitioned is provided by the `CREATE TABLE` statement shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
);
```

(This is the same table used as a basis for the examples in Section 22.2.1, “RANGE Partitioning”. As with the other partitioning examples, we assume that the `default_storage_engine` is `InnoDB`.)

Suppose that there are 20 video stores distributed among 4 franchises as shown in the following table.

Region	Store ID Numbers
North	3, 5, 6, 9, 17
East	1, 2, 10, 11, 19, 20
West	4, 12, 13, 14, 18

Region	Store ID Numbers
Central	7, 8, 15, 16

To partition this table in such a way that rows for stores belonging to the same region are stored in the same partition, you could use the `CREATE TABLE` statement shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
)
PARTITION BY LIST(store_id) (
    PARTITION pNorth VALUES IN (3,5,6,9,17),
    PARTITION pEast VALUES IN (1,2,10,11,19,20),
    PARTITION pWest VALUES IN (4,12,13,14,18),
    PARTITION pCentral VALUES IN (7,8,15,16)
);
```

This makes it easy to add or drop employee records relating to specific regions to or from the table. For instance, suppose that all stores in the West region are sold to another company. In MySQL 8.0, all rows relating to employees working at stores in that region can be deleted with the query `ALTER TABLE employees TRUNCATE PARTITION pWest`, which can be executed much more efficiently than the equivalent `DELETE` statement `DELETE FROM employees WHERE store_id IN (4,12,13,14,18);`. (Using `ALTER TABLE employees DROP PARTITION pWest` would also delete all of these rows, but would also remove the partition `pWest` from the definition of the table; you would need to use an `ALTER TABLE ... ADD PARTITION` statement to restore the table's original partitioning scheme.)

As with `RANGE` partitioning, it is possible to combine `LIST` partitioning with partitioning by hash or key to produce a composite partitioning (subpartitioning). See Section 22.2.6, “Subpartitioning”.

Unlike the case with `RANGE` partitioning, there is no “catch-all” such as `MAXVALUE`; all expected values for the partitioning expression should be covered in `PARTITION ... VALUES IN (...)` clauses. An `INSERT` statement containing an unmatched partitioning column value fails with an error, as shown in this example:

```
mysql> CREATE TABLE h2 (
    ->   c1 INT,
    ->   c2 INT
    -> )
    -> PARTITION BY LIST(c1) (
    ->   PARTITION p0 VALUES IN (1, 4, 7),
    ->   PARTITION p1 VALUES IN (2, 5, 8)
    -> );
Query OK, 0 rows affected (0.11 sec)

mysql> INSERT INTO h2 VALUES (3, 5);
ERROR 1525 (HY000): Table has no partition for value 3
```

When inserting multiple rows using a single `INSERT` statement into a single `InnoDB` table, `InnoDB` considers the statement a single transaction, so that the presence of any unmatched values causes the statement to fail completely, and so no rows are inserted.

You can cause this type of error to be ignored by using the `IGNORE` keyword. If you do so, rows containing unmatched partitioning column values are not inserted, but any rows with matching values *are* inserted, and no errors are reported:

```
mysql> TRUNCATE h2;
Query OK, 1 row affected (0.00 sec)

mysql> SELECT * FROM h2;
Empty set (0.00 sec)

mysql> INSERT IGNORE INTO h2 VALUES (2, 5), (6, 10), (7, 5), (3, 1), (1, 9);
Query OK, 3 rows affected (0.00 sec)
Records: 5  Duplicates: 2  Warnings: 0

mysql> SELECT * FROM h2;
+------+------+
| c1   | c2   |
+------+------+
|    7 |    5 |
|    1 |    9 |
|    2 |    5 |
+------+------+
3 rows in set (0.00 sec)
```

MySQL 8.0 also provides support for `LIST COLUMNS` partitioning, a variant of `LIST` partitioning that enables you to use columns of types other than integer types for partitioning columns, and to use multiple columns as partitioning keys. For more information, see Section 22.2.3.2, “LIST COLUMNS partitioning”.

22.2.3 COLUMNS Partitioning

The next two sections discuss *`COLUMNS` partitioning*, which are variants on `RANGE` and `LIST` partitioning. `COLUMNS` partitioning enables the use of multiple columns in partitioning keys. All of these columns are taken into account both for the purpose of placing rows in partitions and for the determination of which partitions are to be checked for matching rows in partition pruning.

In addition, both `RANGE COLUMNS` partitioning and `LIST COLUMNS` partitioning support the use of non-integer columns for defining value ranges or list members. The permitted data types are shown in the following list:

- All integer types: `TINYINT`, `SMALLINT`, `MEDIUMINT`, `INT` (`INTEGER`), and `BIGINT`. (This is the same as with partitioning by `RANGE` and `LIST`.)

  Other numeric data types (such as `DECIMAL` or `FLOAT`) are not supported as partitioning columns.

- `DATE` and `DATETIME`.

  Columns using other data types relating to dates or times are not supported as partitioning columns.

- The following string types: `CHAR`, `VARCHAR`, `BINARY`, and `VARBINARY`.

  `TEXT` and `BLOB` columns are not supported as partitioning columns.

The discussions of `RANGE COLUMNS` and `LIST COLUMNS` partitioning in the next two sections assume that you are already familiar with partitioning based on ranges and lists as supported in MySQL 5.1 and later; for more information about these, see Section 22.2.1, “RANGE Partitioning”, and Section 22.2.2, “LIST Partitioning”, respectively.

22.2.3.1 RANGE COLUMNS partitioning

Range columns partitioning is similar to range partitioning, but enables you to define partitions using ranges based on multiple column values. In addition, you can define the ranges using columns of types other than integer types.

`RANGE COLUMNS` partitioning differs significantly from `RANGE` partitioning in the following ways:

- `RANGE COLUMNS` does not accept expressions, only names of columns.

- `RANGE COLUMNS` accepts a list of one or more columns.

  `RANGE COLUMNS` partitions are based on comparisons between *tuples* (lists of column values) rather than comparisons between scalar values. Placement of rows in `RANGE COLUMNS` partitions is also based on comparisons between tuples; this is discussed further later in this section.

- `RANGE COLUMNS` partitioning columns are not restricted to integer columns; string, `DATE` and `DATETIME` columns can also be used as partitioning columns. (See Section 22.2.3, "COLUMNS Partitioning", for details.)

The basic syntax for creating a table partitioned by `RANGE COLUMNS` is shown here:

```
CREATE TABLE table_name
PARTITIONED BY RANGE COLUMNS(column_list) (
    PARTITION partition_name VALUES LESS THAN (value_list)[,
    PARTITION partition_name VALUES LESS THAN (value_list)][,
    ...]
)

column_list:
    column_name[, column_name][, ...]

value_list:
    value[, value][, ...]
```

Note

Not all `CREATE TABLE` options that can be used when creating partitioned tables are shown here. For complete information, see Section 13.1.15, "CREATE TABLE Syntax".

In the syntax just shown, *`column_list`* is a list of one or more columns (sometimes called a *partitioning column list*), and *`value_list`* is a list of values (that is, it is a *partition definition value list*). A *`value_list`* must be supplied for each partition definition, and each *`value_list`* must have the same number of values as the *`column_list`* has columns. Generally speaking, if you use *`N`* columns in the `COLUMNS` clause, then each `VALUES LESS THAN` clause must also be supplied with a list of *`N`* values.

The elements in the partitioning column list and in the value list defining each partition must occur in the same order. In addition, each element in the value list must be of the same data type as the corresponding element in the column list. However, the order of the column names in the partitioning column list and the value lists does not have to be the same as the order of the table column definitions in the main part of the `CREATE TABLE` statement. As with table partitioned by `RANGE`, you can use `MAXVALUE` to represent a value such that any legal value inserted into a given column is always less than this value. Here is an example of a `CREATE TABLE` statement that helps to illustrate all of these points:

```
mysql> CREATE TABLE rcx (
    ->     a INT,
    ->     b INT,
    ->     c CHAR(3),
    ->     d INT
    -> )
    -> PARTITION BY RANGE COLUMNS(a,d,c) (
    ->     PARTITION p0 VALUES LESS THAN (5,10,'ggg'),
    ->     PARTITION p1 VALUES LESS THAN (10,20,'mmm'),
    ->     PARTITION p2 VALUES LESS THAN (15,30,'sss'),
    ->     PARTITION p3 VALUES LESS THAN (MAXVALUE,MAXVALUE,MAXVALUE)
    -> );
Query OK, 0 rows affected (0.15 sec)
```

Table `rcx` contains the columns `a`, `b`, `c`, `d`. The partitioning column list supplied to the `COLUMNS` clause uses 3 of these columns, in the order `a`, `d`, `c`. Each value list used to define a partition contains

3 values in the same order; that is, each value list tuple has the form (`INT`, `INT`, `CHAR(3)`), which corresponds to the data types used by columns `a`, `d`, and `c` (in that order).

Placement of rows into partitions is determined by comparing the tuple from a row to be inserted that matches the column list in the `COLUMNS` clause with the tuples used in the `VALUES LESS THAN` clauses to define partitions of the table. Because we are comparing tuples (that is, lists or sets of values) rather than scalar values, the semantics of `VALUES LESS THAN` as used with `RANGE COLUMNS` partitions differs somewhat from the case with simple `RANGE` partitions. In `RANGE` partitioning, a row generating an expression value that is equal to a limiting value in a `VALUES LESS THAN` is never placed in the corresponding partition; however, when using `RANGE COLUMNS` partitioning, it is sometimes possible for a row whose partitioning column list's first element is equal in value to the that of the first element in a `VALUES LESS THAN` value list to be placed in the corresponding partition.

Consider the `RANGE` partitioned table created by this statement:

```
CREATE TABLE r1 (
    a INT,
    b INT
)
PARTITION BY RANGE (a)  (
    PARTITION p0 VALUES LESS THAN (5),
    PARTITION p1 VALUES LESS THAN (MAXVALUE)
);
```

If we insert 3 rows into this table such that the column value for `a` is `5` for each row, all 3 rows are stored in partition `p1` because the `a` column value is in each case not less than 5, as we can see by executing the proper query against the `INFORMATION_SCHEMA.PARTITIONS` table:

```
mysql> INSERT INTO r1 VALUES (5,10), (5,11), (5,12);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT PARTITION_NAME,TABLE_ROWS
    ->     FROM INFORMATION_SCHEMA.PARTITIONS
    ->     WHERE TABLE_NAME = 'r1';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          0 |
| p1             |          3 |
+----------------+------------+
2 rows in set (0.00 sec)
```

Now consider a similar table `rc1` that uses `RANGE COLUMNS` partitioning with both columns `a` and `b` referenced in the `COLUMNS` clause, created as shown here:

```
CREATE TABLE rc1 (
    a INT,
    b INT
)
PARTITION BY RANGE COLUMNS(a, b) (
    PARTITION p0 VALUES LESS THAN (5, 12),
    PARTITION p3 VALUES LESS THAN (MAXVALUE, MAXVALUE)
);
```

If we insert exactly the same rows into `rc1` as we just inserted into `r1`, the distribution of the rows is quite different:

```
mysql> INSERT INTO rc1 VALUES (5,10), (5,11), (5,12);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT PARTITION_NAME,TABLE_ROWS
```

```
    ->     FROM INFORMATION_SCHEMA.PARTITIONS
    ->     WHERE TABLE_NAME = 'rc1';
+--------------+----------------+------------+
| TABLE_SCHEMA | PARTITION_NAME | TABLE_ROWS |
+--------------+----------------+------------+
| p            | p0             |          2 |
| p            | p1             |          1 |
+--------------+----------------+------------+
2 rows in set (0.00 sec)
```

This is because we are comparing rows rather than scalar values. We can compare the row values inserted with the limiting row value from the `VALUES THAN LESS THAN` clause used to define partition `p0` in table `rc1`, like this:

```
mysql> SELECT (5,10) < (5,12), (5,11) < (5,12), (5,12) < (5,12);
+----------------+----------------+----------------+
| (5,10) < (5,12) | (5,11) < (5,12) | (5,12) < (5,12) |
+----------------+----------------+----------------+
|              1 |              1 |              0 |
+----------------+----------------+----------------+
1 row in set (0.00 sec)
```

The 2 tuples `(5,10)` and `(5,11)` evaluate as less than `(5,12)`, so they are stored in partition `p0`. Since 5 is not less than 5 and 12 is not less than 12, `(5,12)` is considered not less than `(5,12)`, and is stored in partition `p1`.

The `SELECT` statement in the preceding example could also have been written using explicit row constructors, like this:

```
SELECT ROW(5,10) < ROW(5,12), ROW(5,11) < ROW(5,12), ROW(5,12) < ROW(5,12);
```

For more information about the use of row constructors in MySQL, see Section 13.2.11.5, “Row Subqueries”.

For a table partitioned by `RANGE COLUMNS` using only a single partitioning column, the storing of rows in partitions is the same as that of an equivalent table that is partitioned by `RANGE`. The following `CREATE TABLE` statement creates a table partitioned by `RANGE COLUMNS` using 1 partitioning column:

```
CREATE TABLE rx (
    a INT,
    b INT
)
PARTITION BY RANGE COLUMNS (a)  (
    PARTITION p0 VALUES LESS THAN (5),
    PARTITION p1 VALUES LESS THAN (MAXVALUE)
);
```

If we insert the rows `(5,10)`, `(5,11)`, and `(5,12)` into this table, we can see that their placement is the same as it is for the table `r` we created and populated earlier:

```
mysql> INSERT INTO rx VALUES (5,10), (5,11), (5,12);
Query OK, 3 rows affected (0.00 sec)
Records: 3  Duplicates: 0  Warnings: 0

mysql> SELECT PARTITION_NAME,TABLE_ROWS
    ->     FROM INFORMATION_SCHEMA.PARTITIONS
    ->     WHERE TABLE_NAME = 'rx';
+--------------+----------------+------------+
| TABLE_SCHEMA | PARTITION_NAME | TABLE_ROWS |
+--------------+----------------+------------+
| p            | p0             |          0 |
| p            | p1             |          3 |
+--------------+----------------+------------+
2 rows in set (0.00 sec)
```

It is also possible to create tables partitioned by `RANGE COLUMNS` where limiting values for one or more columns are repeated in successive partition definitions. You can do this as long as the tuples of column values used to define the partitions are strictly increasing. For example, each of the following `CREATE TABLE` statements is valid:

```
CREATE TABLE rc2 (
    a INT,
    b INT
)
PARTITION BY RANGE COLUMNS(a,b) (
    PARTITION p0 VALUES LESS THAN (0,10),
    PARTITION p1 VALUES LESS THAN (10,20),
    PARTITION p2 VALUES LESS THAN (10,30),
    PARTITION p3 VALUES LESS THAN (MAXVALUE,MAXVALUE)
 );

CREATE TABLE rc3 (
    a INT,
    b INT
)
PARTITION BY RANGE COLUMNS(a,b) (
    PARTITION p0 VALUES LESS THAN (0,10),
    PARTITION p1 VALUES LESS THAN (10,20),
    PARTITION p2 VALUES LESS THAN (10,30),
    PARTITION p3 VALUES LESS THAN (10,35),
    PARTITION p4 VALUES LESS THAN (20,40),
    PARTITION p5 VALUES LESS THAN (MAXVALUE,MAXVALUE)
 );
```

The following statement also succeeds, even though it might appear at first glance that it would not, since the limiting value of column `b` is 25 for partition `p0` and 20 for partition `p1`, and the limiting value of column `c` is 100 for partition `p1` and 50 for partition `p2`:

```
CREATE TABLE rc4 (
    a INT,
    b INT,
    c INT
)
PARTITION BY RANGE COLUMNS(a,b,c) (
    PARTITION p0 VALUES LESS THAN (0,25,50),
    PARTITION p1 VALUES LESS THAN (10,20,100),
    PARTITION p2 VALUES LESS THAN (10,30,50)
    PARTITION p3 VALUES LESS THAN (MAXVALUE,MAXVALUE,MAXVALUE)
 );
```

When designing tables partitioned by `RANGE COLUMNS`, you can always test successive partition definitions by comparing the desired tuples using the `mysql` client, like this:

```
mysql> SELECT (0,25,50) < (10,20,100), (10,20,100) < (10,30,50);
+------------------------+------------------------+
| (0,25,50) < (10,20,100) | (10,20,100) < (10,30,50) |
+------------------------+------------------------+
|                      1 |                      1 |
+------------------------+------------------------+
1 row in set (0.00 sec)
```

If a `CREATE TABLE` statement contains partition definitions that are not in strictly increasing order, it fails with an error, as shown in this example:

```
mysql> CREATE TABLE rcf (
    ->     a INT,
    ->     b INT,
    ->     c INT
    -> )
    -> PARTITION BY RANGE COLUMNS(a,b,c) (
```

```
    ->     PARTITION p0 VALUES LESS THAN (0,25,50),
    ->     PARTITION p1 VALUES LESS THAN (20,20,100),
    ->     PARTITION p2 VALUES LESS THAN (10,30,50),
    ->     PARTITION p3 VALUES LESS THAN (MAXVALUE,MAXVALUE,MAXVALUE)
    ->  );
ERROR 1493 (HY000): VALUES LESS THAN value must be strictly increasing for each partition
```

When you get such an error, you can deduce which partition definitions are invalid by making “less than” comparisons between their column lists. In this case, the problem is with the definition of partition `p2` because the tuple used to define it is not less than the tuple used to define partition `p3`, as shown here:

```
mysql> SELECT (0,25,50) < (20,20,100), (20,20,100) < (10,30,50);
+-------------------------+-------------------------+
| (0,25,50) < (20,20,100) | (20,20,100) < (10,30,50) |
+-------------------------+-------------------------+
|                       1 |                       0 |
+-------------------------+-------------------------+
1 row in set (0.00 sec)
```

It is also possible for `MAXVALUE` to appear for the same column in more than one `VALUES LESS THAN` clause when using `RANGE COLUMNS`. However, the limiting values for individual columns in successive partition definitions should otherwise be increasing, there should be no more than one partition defined where `MAXVALUE` is used as the upper limit for all column values, and this partition definition should appear last in the list of `PARTITION ... VALUES LESS THAN` clauses. In addition, you cannot use `MAXVALUE` as the limiting value for the first column in more than one partition definition.

As stated previously, it is also possible with `RANGE COLUMNS` partitioning to use non-integer columns as partitioning columns. (See Section 22.2.3, “COLUMNS Partitioning”, for a complete listing of these.) Consider a table named `employees` (which is not partitioned), created using the following statement:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
);
```

Using `RANGE COLUMNS` partitioning, you can create a version of this table that stores each row in one of four partitions based on the employee's last name, like this:

```
CREATE TABLE employees_by_lname (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT NOT NULL,
    store_id INT NOT NULL
)
PARTITION BY RANGE COLUMNS (lname)  (
    PARTITION p0 VALUES LESS THAN ('g'),
    PARTITION p1 VALUES LESS THAN ('m'),
    PARTITION p2 VALUES LESS THAN ('t'),
    PARTITION p3 VALUES LESS THAN (MAXVALUE)
);
```

Alternatively, you could cause the `employees` table as created previously to be partitioned using this scheme by executing the following `ALTER TABLE` statement:

```
ALTER TABLE employees PARTITION BY RANGE COLUMNS (lname)  (
```

```
    PARTITION p0 VALUES LESS THAN ('g'),
    PARTITION p1 VALUES LESS THAN ('m'),
    PARTITION p2 VALUES LESS THAN ('t'),
    PARTITION p3 VALUES LESS THAN (MAXVALUE)
);
```

Note

Because different character sets and collations have different sort orders, the character sets and collations in use may effect which partition of a table partitioned by `RANGE COLUMNS` a given row is stored in when using string columns as partitioning columns. In addition, changing the character set or collation for a given database, table, or column after such a table is created may cause changes in how rows are distributed. For example, when using a case-sensitive collation, `'and'` sorts before `'Andersen'`, but when using a collation that is case insensitive, the reverse is true.

For information about how MySQL handles character sets and collations, see Section 10.1, “Character Set Support”.

Similarly, you can cause the `employees` table to be partitioned in such a way that each row is stored in one of several partitions based on the decade in which the corresponding employee was hired using the `ALTER TABLE` statement shown here:

```
ALTER TABLE employees PARTITION BY RANGE COLUMNS (hired)  (
    PARTITION p0 VALUES LESS THAN ('1970-01-01'),
    PARTITION p1 VALUES LESS THAN ('1980-01-01'),
    PARTITION p2 VALUES LESS THAN ('1990-01-01'),
    PARTITION p3 VALUES LESS THAN ('2000-01-01'),
    PARTITION p4 VALUES LESS THAN ('2010-01-01'),
    PARTITION p5 VALUES LESS THAN (MAXVALUE)
);
```

See Section 13.1.15, “CREATE TABLE Syntax”, for additional information about `PARTITION BY RANGE COLUMNS` syntax.

22.2.3.2 LIST COLUMNS partitioning

MySQL 8.0 provides support for `LIST COLUMNS` partitioning. This is a variant of `LIST` partitioning that enables the use of multiple columns as partition keys, and for columns of data types other than integer types to be used as partitioning columns; you can use string types, `DATE`, and `DATETIME` columns. (For more information about permitted data types for `COLUMNS` partitioning columns, see Section 22.2.3, “COLUMNS Partitioning”.)

Suppose that you have a business that has customers in 12 cities which, for sales and marketing purposes, you organize into 4 regions of 3 cities each as shown in the following table:

Region	Cities
1	Oskarshamn, Högsby, Mönsterås
2	Vimmerby, Hultsfred, Västervik
3	Nässjö, Eksjö, Vetlanda
4	Uppvidinge, Alvesta, Växjo

With `LIST COLUMNS` partitioning, you can create a table for customer data that assigns a row to any of 4 partitions corresponding to these regions based on the name of the city where a customer resides, as shown here:

```
CREATE TABLE customers_1 (
```

```
    first_name VARCHAR(25),
    last_name VARCHAR(25),
    street_1 VARCHAR(30),
    street_2 VARCHAR(30),
    city VARCHAR(15),
    renewal DATE
)
PARTITION BY LIST COLUMNS(city) (
    PARTITION pRegion_1 VALUES IN('Oskarshamn', 'Högsby', 'Mönsterås'),
    PARTITION pRegion_2 VALUES IN('Vimmerby', 'Hultsfred', 'Västervik'),
    PARTITION pRegion_3 VALUES IN('Nässjö', 'Eksjö', 'Vetlanda'),
    PARTITION pRegion_4 VALUES IN('Uppvidinge', 'Alvesta', 'Växjo')
);
```

As with partitioning by `RANGE COLUMNS`, you do not need to use expressions in the `COLUMNS()` clause to convert column values into integers. (In fact, the use of expressions other than column names is not permitted with `COLUMNS()`.)

It is also possible to use `DATE` and `DATETIME` columns, as shown in the following example that uses the same name and columns as the `customers_1` table shown previously, but employs `LIST COLUMNS` partitioning based on the `renewal` column to store rows in one of 4 partitions depending on the week in February 2010 the customer's account is scheduled to renew:

```
CREATE TABLE customers_2 (
    first_name VARCHAR(25),
    last_name VARCHAR(25),
    street_1 VARCHAR(30),
    street_2 VARCHAR(30),
    city VARCHAR(15),
    renewal DATE
)
PARTITION BY LIST COLUMNS(renewal) (
    PARTITION pWeek_1 VALUES IN('2010-02-01', '2010-02-02', '2010-02-03',
        '2010-02-04', '2010-02-05', '2010-02-06', '2010-02-07'),
    PARTITION pWeek_2 VALUES IN('2010-02-08', '2010-02-09', '2010-02-10',
        '2010-02-11', '2010-02-12', '2010-02-13', '2010-02-14'),
    PARTITION pWeek_3 VALUES IN('2010-02-15', '2010-02-16', '2010-02-17',
        '2010-02-18', '2010-02-19', '2010-02-20', '2010-02-21'),
    PARTITION pWeek_4 VALUES IN('2010-02-22', '2010-02-23', '2010-02-24',
        '2010-02-25', '2010-02-26', '2010-02-27', '2010-02-28')
);
```

This works, but becomes cumbersome to define and maintain if the number of dates involved grows very large; in such cases, it is usually more practical to employ `RANGE` or `RANGE COLUMNS` partitioning instead. In this case, since the column we wish to use as the partitioning key is a `DATE` column, we use `RANGE COLUMNS` partitioning, as shown here:

```
CREATE TABLE customers_3 (
    first_name VARCHAR(25),
    last_name VARCHAR(25),
    street_1 VARCHAR(30),
    street_2 VARCHAR(30),
    city VARCHAR(15),
    renewal DATE
)
PARTITION BY RANGE COLUMNS(renewal) (
    PARTITION pWeek_1 VALUES LESS THAN('2010-02-09'),
    PARTITION pWeek_2 VALUES LESS THAN('2010-02-15'),
    PARTITION pWeek_3 VALUES LESS THAN('2010-02-22'),
    PARTITION pWeek_4 VALUES LESS THAN('2010-03-01')
);
```

See Section 22.2.3.1, “RANGE COLUMNS partitioning”, for more information.

In addition (as with `RANGE COLUMNS` partitioning), you can use multiple columns in the `COLUMNS()` clause.

See Section 13.1.15, “CREATE TABLE Syntax”, for additional information about `PARTITION BY LIST COLUMNS()` syntax.

22.2.4 HASH Partitioning

Partitioning by `HASH` is used primarily to ensure an even distribution of data among a predetermined number of partitions. With range or list partitioning, you must specify explicitly which partition a given column value or set of column values should be stored in; with hash partitioning, this decision is taken care of for you, and you need only specify a column value or expression based on a column value to be hashed and the number of partitions into which the partitioned table is to be divided.

To partition a table using `HASH` partitioning, it is necessary to append to the `CREATE TABLE` statement a `PARTITION BY HASH (expr)` clause, where *`expr`* is an expression that returns an integer. This can simply be the name of a column whose type is one of MySQL's integer types. In addition, you most likely want to follow this with `PARTITIONS num`, where *`num`* is a positive integer representing the number of partitions into which the table is to be divided.

Note

For simplicity, the tables in the examples that follow do not use any keys. You should be aware that, if a table has any unique keys, every column used in the partitioning expression for this table must be part of every unique key, including the primary key. See Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”, for more information.

The following statement creates a table that uses hashing on the `store_id` column and is divided into 4 partitions:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
)
PARTITION BY HASH(store_id)
PARTITIONS 4;
```

If you do not include a `PARTITIONS` clause, the number of partitions defaults to `1`; using the `PARTITIONS` keyword without a number following it results in a syntax error.

You can also use an SQL expression that returns an integer for *`expr`*. For instance, you might want to partition based on the year in which an employee was hired. This can be done as shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
)
PARTITION BY HASH( YEAR(hired) )
PARTITIONS 4;
```

`expr` must return a nonconstant, nonrandom integer value (in other words, it should be varying but deterministic), and must not contain any prohibited constructs as described in Section 22.6, “Restrictions and Limitations on Partitioning”. You should also keep in mind that this expression is evaluated each time a row is inserted or updated (or possibly deleted); this means that very complex

expressions may give rise to performance issues, particularly when performing operations (such as batch inserts) that affect a great many rows at one time.

The most efficient hashing function is one which operates upon a single table column and whose value increases or decreases consistently with the column value, as this allows for “pruning” on ranges of partitions. That is, the more closely that the expression varies with the value of the column on which it is based, the more efficiently MySQL can use the expression for hash partitioning.

For example, where `date_col` is a column of type `DATE`, then the expression `TO_DAYS(date_col)` is said to vary directly with the value of `date_col`, because for every change in the value of `date_col`, the value of the expression changes in a consistent manner. The variance of the expression `YEAR(date_col)` with respect to `date_col` is not quite as direct as that of `TO_DAYS(date_col)`, because not every possible change in `date_col` produces an equivalent change in `YEAR(date_col)`. Even so, `YEAR(date_col)` is a good candidate for a hashing function, because it varies directly with a portion of `date_col` and there is no possible change in `date_col` that produces a disproportionate change in `YEAR(date_col)`.

By way of contrast, suppose that you have a column named `int_col` whose type is `INT`. Now consider the expression `POW(5-int_col,3) + 6`. This would be a poor choice for a hashing function because a change in the value of `int_col` is not guaranteed to produce a proportional change in the value of the expression. Changing the value of `int_col` by a given amount can produce widely differing changes in the value of the expression. For example, changing `int_col` from `5` to `6` produces a change of `-1` in the value of the expression, but changing the value of `int_col` from `6` to `7` produces a change of `-7` in the expression value.

In other words, the more closely the graph of the column value versus the value of the expression follows a straight line as traced by the equation `y=cx` where `c` is some nonzero constant, the better the expression is suited to hashing. This has to do with the fact that the more nonlinear an expression is, the more uneven the distribution of data among the partitions it tends to produce.

In theory, pruning is also possible for expressions involving more than one column value, but determining which of such expressions are suitable can be quite difficult and time-consuming. For this reason, the use of hashing expressions involving multiple columns is not particularly recommended.

When `PARTITION BY HASH` is used, the storage engine determines which partition of *`num`* partitions to use based on the modulus of the result of the expression. In other words, for a given expression *`expr`*, the partition in which the record is stored is partition number *`N`*, where *`N = MOD(expr, num)`*. Suppose that table `t1` is defined as follows, so that it has 4 partitions:

```
CREATE TABLE t1 (col1 INT, col2 CHAR(5), col3 DATE)
    PARTITION BY HASH( YEAR(col3) )
    PARTITIONS 4;
```

If you insert a record into `t1` whose `col3` value is `'2005-09-15'`, then the partition in which it is stored is determined as follows:

```
MOD(YEAR('2005-09-01'),4)
=  MOD(2005,4)
=  1
```

MySQL 8.0 also supports a variant of `HASH` partitioning known as *linear hashing* which employs a more complex algorithm for determining the placement of new rows inserted into the partitioned table. See Section 22.2.4.1, “LINEAR HASH Partitioning”, for a description of this algorithm.

The user-supplied expression is evaluated each time a record is inserted or updated. It may also—depending on the circumstances—be evaluated when records are deleted.

22.2.4.1 LINEAR HASH Partitioning

MySQL also supports linear hashing, which differs from regular hashing in that linear hashing utilizes a linear powers-of-two algorithm whereas regular hashing employs the modulus of the hashing function's value.

Syntactically, the only difference between linear-hash partitioning and regular hashing is the addition of the `LINEAR` keyword in the `PARTITION BY` clause, as shown here:

```
CREATE TABLE employees (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30),
    hired DATE NOT NULL DEFAULT '1970-01-01',
    separated DATE NOT NULL DEFAULT '9999-12-31',
    job_code INT,
    store_id INT
)
PARTITION BY LINEAR HASH( YEAR(hired) )
PARTITIONS 4;
```

Given an expression *`expr`*, the partition in which the record is stored when linear hashing is used is partition number *`N`* from among *`num`* partitions, where *`N`* is derived according to the following algorithm:

1. Find the next power of 2 greater than *`num`*. We call this value *`V`*; it can be calculated as:

   ```
   V = POWER(2, CEILING(LOG(2, num)))
   ```

   (Suppose that *`num`* is 13. Then `LOG(2,13)` is 3.7004397181411. `CEILING(3.7004397181411)` is 4, and *`V`* = `POWER(2,4)`, which is 16.)

2. Set *`N`* = *`F`*(*`column_list`*) & (*`V`* - 1).

3. While *`N`* >= *`num`*:

   - Set *`V`* = *`V`* / 2
   - Set *`N`* = *`N`* & (*`V`* - 1)

Suppose that the table `t1`, using linear hash partitioning and having 6 partitions, is created using this statement:

```
CREATE TABLE t1 (col1 INT, col2 CHAR(5), col3 DATE)
    PARTITION BY LINEAR HASH( YEAR(col3) )
    PARTITIONS 6;
```

Now assume that you want to insert two records into `t1` having the `col3` column values `'2003-04-14'` and `'1998-10-19'`. The partition number for the first of these is determined as follows:

```
V = POWER(2, CEILING( LOG(2,6) )) = 8
N = YEAR('2003-04-14') & (8 - 1)
   = 2003 & 7
   = 3

(3 >= 6 is FALSE: record stored in partition #3)
```

The number of the partition where the second record is stored is calculated as shown here:

```
V = 8
N = YEAR('1998-10-19') & (8 - 1)
  = 1998 & 7
  = 6
```

```
(6 >= 6 is TRUE: additional step required)

N = 6 & ((8 / 2) - 1)
  = 6 & 3
  = 2

(2 >= 6 is FALSE: record stored in partition #2)
```

The advantage in partitioning by linear hash is that the adding, dropping, merging, and splitting of partitions is made much faster, which can be beneficial when dealing with tables containing extremely large amounts (terabytes) of data. The disadvantage is that data is less likely to be evenly distributed between partitions as compared with the distribution obtained using regular hash partitioning.

22.2.5 KEY Partitioning

Partitioning by key is similar to partitioning by hash, except that where hash partitioning employs a user-defined expression, the hashing function for key partitioning is supplied by the MySQL server. This internal hashing function is based on the same algorithm as `PASSWORD()`.

The syntax rules for `CREATE TABLE ... PARTITION BY KEY` are similar to those for creating a table that is partitioned by hash. The major differences are listed here:

- `KEY` is used rather than `HASH`.
- `KEY` takes only a list of zero or more column names. Any columns used as the partitioning key must comprise part or all of the table's primary key, if the table has one. Where no column name is specified as the partitioning key, the table's primary key is used, if there is one. For example, the following `CREATE TABLE` statement is valid in MySQL 8.0:

  ```
  CREATE TABLE k1 (
      id INT NOT NULL PRIMARY KEY,
      name VARCHAR(20)
  )
  PARTITION BY KEY()
  PARTITIONS 2;
  ```

  If there is no primary key but there is a unique key, then the unique key is used for the partitioning key:

  ```
  CREATE TABLE k1 (
      id INT NOT NULL,
      name VARCHAR(20),
      UNIQUE KEY (id)
  )
  PARTITION BY KEY()
  PARTITIONS 2;
  ```

  However, if the unique key column were not defined as `NOT NULL`, then the previous statement would fail.

  In both of these cases, the partitioning key is the `id` column, even though it is not shown in the output of `SHOW CREATE TABLE` or in the `PARTITION_EXPRESSION` column of the `INFORMATION_SCHEMA.PARTITIONS` table.

  Unlike the case with other partitioning types, columns used for partitioning by `KEY` are not restricted to integer or `NULL` values. For example, the following `CREATE TABLE` statement is valid:

  ```
  CREATE TABLE tm1 (
      s1 CHAR(32) PRIMARY KEY
  )
  PARTITION BY KEY(s1)
  PARTITIONS 10;
  ```

The preceding statement would *not* be valid, were a different partitioning type to be specified. (In this case, simply using `PARTITION BY KEY()` would also be valid and have the same effect as `PARTITION BY KEY(s1)`, since `s1` is the table's primary key.)

For additional information about this issue, see Section 22.6, "Restrictions and Limitations on Partitioning".

Important

For a key-partitioned table, you cannot execute an `ALTER TABLE DROP PRIMARY KEY`, as doing so generates the error `ERROR 1466 (HY000): Field in list of fields for partition function not found in table`.

It is also possible to partition a table by linear key. Here is a simple example:

```
CREATE TABLE tk (
    col1 INT NOT NULL,
    col2 CHAR(5),
    col3 DATE
)
PARTITION BY LINEAR KEY (col1)
PARTITIONS 3;
```

The `LINEAR` keyword has the same effect on `KEY` partitioning as it does on `HASH` partitioning, with the partition number being derived using a powers-of-two algorithm rather than modulo arithmetic. See Section 22.2.4.1, "LINEAR HASH Partitioning", for a description of this algorithm and its implications.

22.2.6 Subpartitioning

Subpartitioning—also known as *composite partitioning*—is the further division of each partition in a partitioned table. Consider the following `CREATE TABLE` statement:

```
CREATE TABLE ts (id INT, purchased DATE)
    PARTITION BY RANGE( YEAR(purchased) )
    SUBPARTITION BY HASH( TO_DAYS(purchased) )
    SUBPARTITIONS 2 (
        PARTITION p0 VALUES LESS THAN (1990),
        PARTITION p1 VALUES LESS THAN (2000),
        PARTITION p2 VALUES LESS THAN MAXVALUE
    );
```

Table `ts` has 3 `RANGE` partitions. Each of these partitions—`p0`, `p1`, and `p2`—is further divided into 2 subpartitions. In effect, the entire table is divided into `3 * 2 = 6` partitions. However, due to the action of the `PARTITION BY RANGE` clause, the first 2 of these store only those records with a value less than 1990 in the `purchased` column.

It is possible to subpartition tables that are partitioned by `RANGE` or `LIST`. Subpartitions may use either `HASH` or `KEY` partitioning. This is also known as *composite partitioning*.

Note

`SUBPARTITION BY HASH` and `SUBPARTITION BY KEY` generally follow the same syntax rules as `PARTITION BY HASH` and `PARTITION BY KEY`, respectively. An exception to this is that `SUBPARTITION BY KEY` (unlike `PARTITION BY KEY`) does not currently support a default column, so the column used for this purpose must be specified, even if the table has an explicit primary key. This is a known issue which we are working to address; see Issues with subpartitions, for more information and an example.

It is also possible to define subpartitions explicitly using `SUBPARTITION` clauses to specify options for individual subpartitions. For example, a more verbose fashion of creating the same table `ts` as shown in the previous example would be:

```
CREATE TABLE ts (id INT, purchased DATE)
    PARTITION BY RANGE( YEAR(purchased) )
    SUBPARTITION BY HASH( TO_DAYS(purchased) ) (
        PARTITION p0 VALUES LESS THAN (1990) (
            SUBPARTITION s0,
            SUBPARTITION s1
        ),
        PARTITION p1 VALUES LESS THAN (2000) (
            SUBPARTITION s2,
            SUBPARTITION s3
        ),
        PARTITION p2 VALUES LESS THAN MAXVALUE (
            SUBPARTITION s4,
            SUBPARTITION s5
        )
    );
```

Some syntactical items of note are listed here:

- Each partition must have the same number of subpartitions.
- If you explicitly define any subpartitions using `SUBPARTITION` on any partition of a partitioned table, you must define them all. In other words, the following statement will fail:

  ```
  CREATE TABLE ts (id INT, purchased DATE)
      PARTITION BY RANGE( YEAR(purchased) )
      SUBPARTITION BY HASH( TO_DAYS(purchased) ) (
          PARTITION p0 VALUES LESS THAN (1990) (
              SUBPARTITION s0,
              SUBPARTITION s1
          ),
          PARTITION p1 VALUES LESS THAN (2000),
          PARTITION p2 VALUES LESS THAN MAXVALUE (
              SUBPARTITION s2,
              SUBPARTITION s3
          )
      );
  ```

  This statement would still fail even if it used `SUBPARTITIONS 2`.

- Each `SUBPARTITION` clause must include (at a minimum) a name for the subpartition. Otherwise, you may set any desired option for the subpartition or allow it to assume its default setting for that option.
- Subpartition names must be unique across the entire table. For example, the following `CREATE TABLE` statement is valid:

  ```
  CREATE TABLE ts (id INT, purchased DATE)
      PARTITION BY RANGE( YEAR(purchased) )
      SUBPARTITION BY HASH( TO_DAYS(purchased) ) (
          PARTITION p0 VALUES LESS THAN (1990) (
              SUBPARTITION s0,
              SUBPARTITION s1
          ),
          PARTITION p1 VALUES LESS THAN (2000) (
              SUBPARTITION s2,
              SUBPARTITION s3
          ),
          PARTITION p2 VALUES LESS THAN MAXVALUE (
              SUBPARTITION s4,
              SUBPARTITION s5
          )
  ```

```
    );
```

22.2.7 How MySQL Partitioning Handles NULL

Partitioning in MySQL does nothing to disallow `NULL` as the value of a partitioning expression, whether it is a column value or the value of a user-supplied expression. Even though it is permitted to use `NULL` as the value of an expression that must otherwise yield an integer, it is important to keep in mind that `NULL` is not a number. MySQL's partitioning implementation treats `NULL` as being less than any non-`NULL` value, just as `ORDER BY` does.

This means that treatment of `NULL` varies between partitioning of different types, and may produce behavior which you do not expect if you are not prepared for it. This being the case, we discuss in this section how each MySQL partitioning type handles `NULL` values when determining the partition in which a row should be stored, and provide examples for each.

Handling of NULL with RANGE partitioning. If you insert a row into a table partitioned by `RANGE` such that the column value used to determine the partition is `NULL`, the row is inserted into the lowest partition. Consider these two tables in a database named `p`, created as follows:

```
mysql> CREATE TABLE t1 (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY RANGE(c1) (
    ->     PARTITION p0 VALUES LESS THAN (0),
    ->     PARTITION p1 VALUES LESS THAN (10),
    ->     PARTITION p2 VALUES LESS THAN MAXVALUE
    -> );
Query OK, 0 rows affected (0.09 sec)

mysql> CREATE TABLE t2 (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY RANGE(c1) (
    ->     PARTITION p0 VALUES LESS THAN (-5),
    ->     PARTITION p1 VALUES LESS THAN (0),
    ->     PARTITION p2 VALUES LESS THAN (10),
    ->     PARTITION p3 VALUES LESS THAN MAXVALUE
    -> );
Query OK, 0 rows affected (0.09 sec)
```

You can see the partitions created by these two `CREATE TABLE` statements using the following query against the `PARTITIONS` table in the `INFORMATION_SCHEMA` database:

```
mysql> SELECT TABLE_NAME, PARTITION_NAME, TABLE_ROWS, AVG_ROW_LENGTH, DATA_LENGTH
     >   FROM INFORMATION_SCHEMA.PARTITIONS
     >   WHERE TABLE_SCHEMA = 'p' AND TABLE_NAME LIKE 't_';
+------------+----------------+------------+----------------+-------------+
| TABLE_NAME | PARTITION_NAME | TABLE_ROWS | AVG_ROW_LENGTH | DATA_LENGTH |
+------------+----------------+------------+----------------+-------------+
| t1         | p0             |          0 |              0 |           0 |
| t1         | p1             |          0 |              0 |           0 |
| t1         | p2             |          0 |              0 |           0 |
| t2         | p0             |          0 |              0 |           0 |
| t2         | p1             |          0 |              0 |           0 |
| t2         | p2             |          0 |              0 |           0 |
| t2         | p3             |          0 |              0 |           0 |
+------------+----------------+------------+----------------+-------------+
7 rows in set (0.00 sec)
```

(For more information about this table, see Section 24.13, “The INFORMATION_SCHEMA PARTITIONS Table”.) Now let us populate each of these tables with a single row containing a `NULL` in the column used as the partitioning key, and verify that the rows were inserted using a pair of `SELECT` statements:

```
mysql> INSERT INTO t1 VALUES (NULL, 'mothra');
Query OK, 1 row affected (0.00 sec)

mysql> INSERT INTO t2 VALUES (NULL, 'mothra');
Query OK, 1 row affected (0.00 sec)

mysql> SELECT * FROM t1;
+------+--------+
| id   | name   |
+------+--------+
| NULL | mothra |
+------+--------+
1 row in set (0.00 sec)

mysql> SELECT * FROM t2;
+------+--------+
| id   | name   |
+------+--------+
| NULL | mothra |
+------+--------+
1 row in set (0.00 sec)
```

You can see which partitions are used to store the inserted rows by rerunning the previous query against `INFORMATION_SCHEMA.PARTITIONS` and inspecting the output:

```
mysql> SELECT TABLE_NAME, PARTITION_NAME, TABLE_ROWS, AVG_ROW_LENGTH, DATA_LENGTH
     >   FROM INFORMATION_SCHEMA.PARTITIONS
     >   WHERE TABLE_SCHEMA = 'p' AND TABLE_NAME LIKE 't_';
+------------+----------------+------------+----------------+-------------+
| TABLE_NAME | PARTITION_NAME | TABLE_ROWS | AVG_ROW_LENGTH | DATA_LENGTH |
+------------+----------------+------------+----------------+-------------+
| t1         | p0             |          1 |             20 |          20 |
| t1         | p1             |          0 |              0 |           0 |
| t1         | p2             |          0 |              0 |           0 |
| t2         | p0             |          1 |             20 |          20 |
| t2         | p1             |          0 |              0 |           0 |
| t2         | p2             |          0 |              0 |           0 |
| t2         | p3             |          0 |              0 |           0 |
+------------+----------------+------------+----------------+-------------+
7 rows in set (0.01 sec)
```

You can also demonstrate that these rows were stored in the lowest-numbered partition of each table by dropping these partitions, and then re-running the `SELECT` statements:

```
mysql> ALTER TABLE t1 DROP PARTITION p0;
Query OK, 0 rows affected (0.16 sec)

mysql> ALTER TABLE t2 DROP PARTITION p0;
Query OK, 0 rows affected (0.16 sec)

mysql> SELECT * FROM t1;
Empty set (0.00 sec)

mysql> SELECT * FROM t2;
Empty set (0.00 sec)
```

(For more information on `ALTER TABLE ... DROP PARTITION`, see Section 13.1.7, “ALTER TABLE Syntax”.)

`NULL` is also treated in this way for partitioning expressions that use SQL functions. Suppose that we define a table using a `CREATE TABLE` statement such as this one:

```
CREATE TABLE tndate (
    id INT,
    dt DATE
)
PARTITION BY RANGE( YEAR(dt) ) (
```

```
    PARTITION p0 VALUES LESS THAN (1990),
    PARTITION p1 VALUES LESS THAN (2000),
    PARTITION p2 VALUES LESS THAN MAXVALUE
);
```

As with other MySQL functions, `YEAR(NULL)` returns `NULL`. A row with a `dt` column value of `NULL` is treated as though the partitioning expression evaluated to a value less than any other value, and so is inserted into partition `p0`.

Handling of NULL with LIST partitioning. A table that is partitioned by `LIST` admits `NULL` values if and only if one of its partitions is defined using that value-list that contains `NULL`. The converse of this is that a table partitioned by `LIST` which does not explicitly use `NULL` in a value list rejects rows resulting in a `NULL` value for the partitioning expression, as shown in this example:

```
mysql> CREATE TABLE ts1 (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY LIST(c1) (
    ->     PARTITION p0 VALUES IN (0, 3, 6),
    ->     PARTITION p1 VALUES IN (1, 4, 7),
    ->     PARTITION p2 VALUES IN (2, 5, 8)
    -> );
Query OK, 0 rows affected (0.01 sec)

mysql> INSERT INTO ts1 VALUES (9, 'mothra');
ERROR 1504 (HY000): Table has no partition for value 9

mysql> INSERT INTO ts1 VALUES (NULL, 'mothra');
ERROR 1504 (HY000): Table has no partition for value NULL
```

Only rows having a `c1` value between `0` and `8` inclusive can be inserted into `ts1`. `NULL` falls outside this range, just like the number `9`. We can create tables `ts2` and `ts3` having value lists containing `NULL`, as shown here:

```
mysql> CREATE TABLE ts2 (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY LIST(c1) (
    ->     PARTITION p0 VALUES IN (0, 3, 6),
    ->     PARTITION p1 VALUES IN (1, 4, 7),
    ->     PARTITION p2 VALUES IN (2, 5, 8),
    ->     PARTITION p3 VALUES IN (NULL)
    -> );
Query OK, 0 rows affected (0.01 sec)

mysql> CREATE TABLE ts3 (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY LIST(c1) (
    ->     PARTITION p0 VALUES IN (0, 3, 6),
    ->     PARTITION p1 VALUES IN (1, 4, 7, NULL),
    ->     PARTITION p2 VALUES IN (2, 5, 8)
    -> );
Query OK, 0 rows affected (0.01 sec)
```

When defining value lists for partitioning, you can (and should) treat `NULL` just as you would any other value. For example, both `VALUES IN (NULL)` and `VALUES IN (1, 4, 7, NULL)` are valid, as are `VALUES IN (1, NULL, 4, 7)`, `VALUES IN (NULL, 1, 4, 7)`, and so on. You can insert a row having `NULL` for column `c1` into each of the tables `ts2` and `ts3`:

```
mysql> INSERT INTO ts2 VALUES (NULL, 'mothra');
Query OK, 1 row affected (0.00 sec)
```

```
mysql> INSERT INTO ts3 VALUES (NULL, 'mothra');
Query OK, 1 row affected (0.00 sec)
```

By issuing the appropriate query against `INFORMATION_SCHEMA.PARTITIONS`, you can determine which partitions were used to store the rows just inserted (we assume, as in the previous examples, that the partitioned tables were created in the `p` database):

```
mysql> SELECT TABLE_NAME, PARTITION_NAME, TABLE_ROWS, AVG_ROW_LENGTH, DATA_LENGTH
     >   FROM INFORMATION_SCHEMA.PARTITIONS
     >   WHERE TABLE_SCHEMA = 'p' AND TABLE_NAME LIKE 'ts_';
+------------+----------------+------------+----------------+-------------+
| TABLE_NAME | PARTITION_NAME | TABLE_ROWS | AVG_ROW_LENGTH | DATA_LENGTH |
+------------+----------------+------------+----------------+-------------+
| ts2        | p0             |          0 |              0 |           0 |
| ts2        | p1             |          0 |              0 |           0 |
| ts2        | p2             |          0 |              0 |           0 |
| ts2        | p3             |          1 |             20 |          20 |
| ts3        | p0             |          0 |              0 |           0 |
| ts3        | p1             |          1 |             20 |          20 |
| ts3        | p2             |          0 |              0 |           0 |
+------------+----------------+------------+----------------+-------------+
7 rows in set (0.01 sec)
```

As shown earlier in this section, you can also verify which partitions were used for storing the rows by deleting these partitions and then performing a `SELECT`.

Handling of NULL with HASH and KEY partitioning. `NULL` is handled somewhat differently for tables partitioned by `HASH` or `KEY`. In these cases, any partition expression that yields a `NULL` value is treated as though its return value were zero. We can verify this behavior by examining the effects on the file system of creating a table partitioned by `HASH` and populating it with a record containing appropriate values. Suppose that you have a table `th` (also in the `p` database) created using the following statement:

```
mysql> CREATE TABLE th (
    ->     c1 INT,
    ->     c2 VARCHAR(20)
    -> )
    -> PARTITION BY HASH(c1)
    -> PARTITIONS 2;
Query OK, 0 rows affected (0.00 sec)
```

The partitions belonging to this table can be viewed using the query shown here:

```
mysql> SELECT TABLE_NAME,PARTITION_NAME,TABLE_ROWS,AVG_ROW_LENGTH,DATA_LENGTH
     >   FROM INFORMATION_SCHEMA.PARTITIONS
     >   WHERE TABLE_SCHEMA = 'p' AND TABLE_NAME ='th';
+------------+----------------+------------+----------------+-------------+
| TABLE_NAME | PARTITION_NAME | TABLE_ROWS | AVG_ROW_LENGTH | DATA_LENGTH |
+------------+----------------+------------+----------------+-------------+
| th         | p0             |          0 |              0 |           0 |
| th         | p1             |          0 |              0 |           0 |
+------------+----------------+------------+----------------+-------------+
2 rows in set (0.00 sec)
```

`TABLE_ROWS` for each partition is 0. Now insert two rows into `th` whose `c1` column values are `NULL` and 0, and verify that these rows were inserted, as shown here:

```
mysql> INSERT INTO th VALUES (NULL, 'mothra'), (0, 'gigan');
Query OK, 1 row affected (0.00 sec)

mysql> SELECT * FROM th;
+------+---------+
| c1   | c2      |
+------+---------+
| NULL | mothra  |
+------+---------+
```

```
|    0 | gigan   |
+------+---------+
2 rows in set (0.01 sec)
```

Recall that for any integer *`N`*, the value of `NULL MOD N` is always `NULL`. For tables that are partitioned by `HASH` or `KEY`, this result is treated for determining the correct partition as `0`. Checking the `INFORMATION_SCHEMA.PARTITIONS` table once again, we can see that both rows were inserted into partition `p0`:

```
mysql> SELECT TABLE_NAME, PARTITION_NAME, TABLE_ROWS, AVG_ROW_LENGTH, DATA_LENGTH
    >   FROM INFORMATION_SCHEMA.PARTITIONS
    >   WHERE TABLE_SCHEMA = 'p' AND TABLE_NAME ='th';
+------------+----------------+------------+----------------+-------------+
| TABLE_NAME | PARTITION_NAME | TABLE_ROWS | AVG_ROW_LENGTH | DATA_LENGTH |
+------------+----------------+------------+----------------+-------------+
| th         | p0             |          2 |             20 |          20 |
| th         | p1             |          0 |              0 |           0 |
+------------+----------------+------------+----------------+-------------+
2 rows in set (0.00 sec)
```

By repeating the last example using `PARTITION BY KEY` in place of `PARTITION BY HASH` in the definition of the table, you can verify that `NULL` is also treated like 0 for this type of partitioning.

22.3 Partition Management

There are a number of ways using SQL statements to modify partitioned tables; it is possible to add, drop, redefine, merge, or split existing partitions using the partitioning extensions to the `ALTER TABLE` statement. There are also ways to obtain information about partitioned tables and partitions. We discuss these topics in the sections that follow.

- For information about partition management in tables partitioned by `RANGE` or `LIST`, see Section 22.3.1, “Management of RANGE and LIST Partitions”.
- For a discussion of managing `HASH` and `KEY` partitions, see Section 22.3.2, “Management of HASH and KEY Partitions”.
- See Section 22.3.5, “Obtaining Information About Partitions”, for a discussion of mechanisms provided in MySQL 8.0 for obtaining information about partitioned tables and partitions.
- For a discussion of performing maintenance operations on partitions, see Section 22.3.4, “Maintenance of Partitions”.

Note

All partitions of a partitioned table must have the same number of subpartitions; it is not possible to change the subpartitioning once the table has been created.

To change a table's partitioning scheme, it is necessary only to use the `ALTER TABLE` statement with a *`partition_options`* option, which has the same syntax as that as used with `CREATE TABLE` for creating a partitioned table; this option (also) always begins with the keywords `PARTITION BY`. Suppose that the following `CREATE TABLE` statement was used to create a table that is partitioned by range:

```
CREATE TABLE trb3 (id INT, name VARCHAR(50), purchased DATE)
    PARTITION BY RANGE( YEAR(purchased) ) (
        PARTITION p0 VALUES LESS THAN (1990),
        PARTITION p1 VALUES LESS THAN (1995),
        PARTITION p2 VALUES LESS THAN (2000),
        PARTITION p3 VALUES LESS THAN (2005)
    );
```

To repartition this table so that it is partitioned by key into two partitions using the `id` column value as the basis for the key, you can use this statement:

```
ALTER TABLE trb3 PARTITION BY KEY(id) PARTITIONS 2;
```

This has the same effect on the structure of the table as dropping the table and re-creating it using `CREATE TABLE trb3 PARTITION BY KEY(id) PARTITIONS 2;`.

`ALTER TABLE ... ENGINE = ...` changes only the storage engine used by the table, and leaves the table's partitioning scheme intact. The statement succeeds only if the target storage engine provides partitioning support. You can use `ALTER TABLE ... REMOVE PARTITIONING` to remove a table's partitioning; see Section 13.1.7, "ALTER TABLE Syntax".

Important

Only a single `PARTITION BY`, `ADD PARTITION`, `DROP PARTITION`, `REORGANIZE PARTITION`, or `COALESCE PARTITION` clause can be used in a given `ALTER TABLE` statement. If you (for example) wish to drop a partition and reorganize a table's remaining partitions, you must do so in two separate `ALTER TABLE` statements (one using `DROP PARTITION` and then a second one using `REORGANIZE PARTITION`).

You can delete all rows from one or more selected partitions using `ALTER TABLE ... TRUNCATE PARTITION`.

22.3.1 Management of RANGE and LIST Partitions

Adding and dropping of range and list partitions are handled in a similar fashion, so we discuss the management of both sorts of partitioning in this section. For information about working with tables that are partitioned by hash or key, see Section 22.3.2, "Management of HASH and KEY Partitions".

Dropping a partition from a table that is partitioned by either `RANGE` or by `LIST` can be accomplished using the `ALTER TABLE` statement with the `DROP PARTITION` option. Suppose that you have created a table that is partitioned by range and then populated with 10 records using the following `CREATE TABLE` and `INSERT` statements:

```
mysql> CREATE TABLE tr (id INT, name VARCHAR(50), purchased DATE)
    ->     PARTITION BY RANGE( YEAR(purchased) ) (
    ->         PARTITION p0 VALUES LESS THAN (1990),
    ->         PARTITION p1 VALUES LESS THAN (1995),
    ->         PARTITION p2 VALUES LESS THAN (2000),
    ->         PARTITION p3 VALUES LESS THAN (2005),
    ->         PARTITION p4 VALUES LESS THAN (2010),
    ->         PARTITION p5 VALUES LESS THAN (2015)
    ->     );
Query OK, 0 rows affected (0.28 sec)

mysql> INSERT INTO tr VALUES
    ->     (1, 'desk organiser', '2003-10-15'),
    ->     (2, 'alarm clock', '1997-11-05'),
    ->     (3, 'chair', '2009-03-10'),
    ->     (4, 'bookcase', '1989-01-10'),
    ->     (5, 'exercise bike', '2014-05-09'),
    ->     (6, 'sofa', '1987-06-05'),
    ->     (7, 'espresso maker', '2011-11-22'),
    ->     (8, 'aquarium', '1992-08-04'),
    ->     (9, 'study desk', '2006-09-16'),
    ->     (10, 'lava lamp', '1998-12-25');
Query OK, 10 rows affected (0.05 sec)
Records: 10  Duplicates: 0  Warnings: 0
```

You can see which items should have been inserted into partition `p2` as shown here:

```
mysql> SELECT * FROM tr
    ->     WHERE purchased BETWEEN '1995-01-01' AND '1999-12-31';
+------+-------------+------------+
```

```
| id   | name        | purchased  |
+------+-------------+------------+
|    2 | alarm clock | 1997-11-05 |
|   10 | lava lamp   | 1998-12-25 |
+------+-------------+------------+
2 rows in set (0.00 sec)
```

You can also get this information using partition selection, as shown here:

```
mysql> SELECT * FROM tr PARTITION (p2);
+------+-------------+------------+
| id   | name        | purchased  |
+------+-------------+------------+
|    2 | alarm clock | 1997-11-05 |
|   10 | lava lamp   | 1998-12-25 |
+------+-------------+------------+
2 rows in set (0.00 sec)
```

See Section 22.5, “Partition Selection”, for more information.

To drop the partition named `p2`, execute the following command:

```
mysql> ALTER TABLE tr DROP PARTITION p2;
Query OK, 0 rows affected (0.03 sec)
```

Note

The `NDBCLUSTER` storage engine does not support `ALTER TABLE ... DROP PARTITION`. It does, however, support the other partitioning-related extensions to `ALTER TABLE` that are described in this chapter.

It is very important to remember that, *when you drop a partition, you also delete all the data that was stored in that partition*. You can see that this is the case by re-running the previous `SELECT` query:

```
mysql> SELECT * FROM tr WHERE purchased
    -> BETWEEN '1995-01-01' AND '1999-12-31';
Empty set (0.00 sec)
```

Note

As of MySQL 8.0.2, `DROP PARTITION` is supported by native partitioning in-place APIs and may be used with `ALGORITHM={COPY|INPLACE}`. `DROP PARTITION` with `ALGORITHM=INPLACE` deletes data stored in the partition and drops the partition. However, `DROP PARTITION` with `ALGORITHM=COPY` or `old_alter_table=ON` rebuilds the partitioned table and attempts to move data from the dropped partition to another partition with a compatible `PARTITION ... VALUES` definition. Data that cannot be moved to another partition is deleted.

Because of this, you must have the `DROP` privilege for a table before you can execute `ALTER TABLE ... DROP PARTITION` on that table.

If you wish to drop all data from all partitions while preserving the table definition and its partitioning scheme, use the `TRUNCATE TABLE` statement. (See Section 13.1.30, “TRUNCATE TABLE Syntax”.)

If you intend to change the partitioning of a table *without* losing data, use `ALTER TABLE ... REORGANIZE PARTITION` instead. See below or in Section 13.1.7, “ALTER TABLE Syntax”, for information about `REORGANIZE PARTITION`.

If you now execute a `SHOW CREATE TABLE` statement, you can see how the partitioning makeup of the table has been changed:

```
mysql> SHOW CREATE TABLE tr\G
```

```
*************************** 1. row ***************************
       Table: tr
Create Table: CREATE TABLE `tr` (
  `id` int(11) DEFAULT NULL,
  `name` varchar(50) DEFAULT NULL,
  `purchased` date DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=latin1
/*!50100 PARTITION BY RANGE ( YEAR(purchased))
(PARTITION p0 VALUES LESS THAN (1990) ENGINE = InnoDB,
 PARTITION p1 VALUES LESS THAN (1995) ENGINE = InnoDB,
 PARTITION p3 VALUES LESS THAN (2005) ENGINE = InnoDB,
 PARTITION p4 VALUES LESS THAN (2010) ENGINE = InnoDB,
 PARTITION p5 VALUES LESS THAN (2015) ENGINE = InnoDB) */
1 row in set (0.00 sec)
```

When you insert new rows into the changed table with `purchased` column values between `'1995-01-01'` and `'2004-12-31'` inclusive, those rows will be stored in partition `p3`. You can verify this as follows:

```
mysql> INSERT INTO tr VALUES (11, 'pencil holder', '1995-07-12');
Query OK, 1 row affected (0.00 sec)

mysql> SELECT * FROM tr WHERE purchased
    -> BETWEEN '1995-01-01' AND '2004-12-31';
+------+---------------+------------+
| id   | name          | purchased  |
+------+---------------+------------+
|    1 | desk organiser | 2003-10-15 |
|   11 | pencil holder  | 1995-07-12 |
+------+---------------+------------+
2 rows in set (0.00 sec)

mysql> ALTER TABLE tr DROP PARTITION p3;
Query OK, 0 rows affected (0.03 sec)

mysql> SELECT * FROM tr WHERE purchased
    -> BETWEEN '1995-01-01' AND '2004-12-31';
Empty set (0.00 sec)
```

The number of rows dropped from the table as a result of `ALTER TABLE ... DROP PARTITION` is not reported by the server as it would be by the equivalent `DELETE` query.

Dropping `LIST` partitions uses exactly the same `ALTER TABLE ... DROP PARTITION` syntax as used for dropping `RANGE` partitions. However, there is one important difference in the effect this has on your use of the table afterward: You can no longer insert into the table any rows having any of the values that were included in the value list defining the deleted partition. (See Section 22.2.2, “LIST Partitioning”, for an example.)

To add a new range or list partition to a previously partitioned table, use the `ALTER TABLE ... ADD PARTITION` statement. For tables which are partitioned by `RANGE`, this can be used to add a new range to the end of the list of existing partitions. Suppose that you have a partitioned table containing membership data for your organization, which is defined as follows:

```
CREATE TABLE members (
    id INT,
    fname VARCHAR(25),
    lname VARCHAR(25),
    dob DATE
)
PARTITION BY RANGE( YEAR(dob) ) (
    PARTITION p0 VALUES LESS THAN (1980),
    PARTITION p1 VALUES LESS THAN (1990),
    PARTITION p2 VALUES LESS THAN (2000)
);
```

Suppose further that the minimum age for members is 16. As the calendar approaches the end of 2015, you realize that you will soon be admitting members who were born in 2000 (and later). You can

modify the `members` table to accommodate new members born in the years 2000 to 2010 as shown here:

```
ALTER TABLE members ADD PARTITION (PARTITION p3 VALUES LESS THAN (2010));
```

With tables that are partitioned by range, you can use `ADD PARTITION` to add new partitions to the high end of the partitions list only. Trying to add a new partition in this manner between or before existing partitions results in an error as shown here:

```
mysql> ALTER TABLE members
    >     ADD PARTITION (
    >     PARTITION n VALUES LESS THAN (1970));
ERROR 1463 (HY000): VALUES LESS THAN value must be strictly »
   increasing for each partition
```

You can work around this problem by reorganizing the first partition into two new ones that split the range between them, like this:

```
ALTER TABLE members
    REORGANIZE PARTITION p0 INTO (
        PARTITION n0 VALUES LESS THAN (1970),
        PARTITION n1 VALUES LESS THAN (1980)
);
```

Using `SHOW CREATE TABLE` you can see that the `ALTER TABLE` statement has had the desired effect:

```
mysql> SHOW CREATE TABLE members\G
*************************** 1. row ***************************
       Table: members
Create Table: CREATE TABLE `members` (
  `id` int(11) DEFAULT NULL,
  `fname` varchar(25) DEFAULT NULL,
  `lname` varchar(25) DEFAULT NULL,
  `dob` date DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=latin1
/*!50100 PARTITION BY RANGE ( YEAR(dob))
(PARTITION n0 VALUES LESS THAN (1970) ENGINE = InnoDB,
 PARTITION n1 VALUES LESS THAN (1980) ENGINE = InnoDB,
 PARTITION p1 VALUES LESS THAN (1990) ENGINE = InnoDB,
 PARTITION p2 VALUES LESS THAN (2000) ENGINE = InnoDB,
 PARTITION p3 VALUES LESS THAN (2010) ENGINE = InnoDB) */
1 row in set (0.00 sec)
```

See also Section 13.1.7.1, “ALTER TABLE Partition Operations”.

You can also use `ALTER TABLE ... ADD PARTITION` to add new partitions to a table that is partitioned by `LIST`. Suppose a table `tt` is defined using the following `CREATE TABLE` statement:

```
CREATE TABLE tt (
    id INT,
    data INT
)
PARTITION BY LIST(data) (
    PARTITION p0 VALUES IN (5, 10, 15),
    PARTITION p1 VALUES IN (6, 12, 18)
);
```

You can add a new partition in which to store rows having the `data` column values `7`, `14`, and `21` as shown:

```
ALTER TABLE tt ADD PARTITION (PARTITION p2 VALUES IN (7, 14, 21));
```

Keep in mind that you *cannot* add a new `LIST` partition encompassing any values that are already included in the value list of an existing partition. If you attempt to do so, an error will result:

```
mysql> ALTER TABLE tt ADD PARTITION
    >     (PARTITION np VALUES IN (4, 8, 12));
ERROR 1465 (HY000): Multiple definition of same constant »
                    in list partitioning
```

Because any rows with the `data` column value `12` have already been assigned to partition `p1`, you cannot create a new partition on table `tt` that includes `12` in its value list. To accomplish this, you could drop `p1`, and add `np` and then a new `p1` with a modified definition. However, as discussed earlier, this would result in the loss of all data stored in `p1`—and it is often the case that this is not what you really want to do. Another solution might appear to be to make a copy of the table with the new partitioning and to copy the data into it using `CREATE TABLE ... SELECT ...`, then drop the old table and rename the new one, but this could be very time-consuming when dealing with a large amounts of data. This also might not be feasible in situations where high availability is a requirement.

You can add multiple partitions in a single `ALTER TABLE ... ADD PARTITION` statement as shown here:

```
CREATE TABLE employees (
  id INT NOT NULL,
  fname VARCHAR(50) NOT NULL,
  lname VARCHAR(50) NOT NULL,
  hired DATE NOT NULL
)
PARTITION BY RANGE( YEAR(hired) ) (
  PARTITION p1 VALUES LESS THAN (1991),
  PARTITION p2 VALUES LESS THAN (1996),
  PARTITION p3 VALUES LESS THAN (2001),
  PARTITION p4 VALUES LESS THAN (2005)
);

ALTER TABLE employees ADD PARTITION (
    PARTITION p5 VALUES LESS THAN (2010),
    PARTITION p6 VALUES LESS THAN MAXVALUE
);
```

Fortunately, MySQL's partitioning implementation provides ways to redefine partitions without losing data. Let us look first at a couple of simple examples involving `RANGE` partitioning. Recall the `members` table which is now defined as shown here:

```
mysql> SHOW CREATE TABLE members\G
*************************** 1. row ***************************
       Table: members
Create Table: CREATE TABLE `members` (
  `id` int(11) DEFAULT NULL,
  `fname` varchar(25) DEFAULT NULL,
  `lname` varchar(25) DEFAULT NULL,
  `dob` date DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=latin1
/*!50100 PARTITION BY RANGE ( YEAR(dob))
(PARTITION n0 VALUES LESS THAN (1970) ENGINE = InnoDB,
 PARTITION n1 VALUES LESS THAN (1980) ENGINE = InnoDB,
 PARTITION p1 VALUES LESS THAN (1990) ENGINE = InnoDB,
 PARTITION p2 VALUES LESS THAN (2000) ENGINE = InnoDB,
 PARTITION p3 VALUES LESS THAN (2010) ENGINE = InnoDB) */
1 row in set (0.00 sec)
```

Suppose that you would like to move all rows representing members born before 1960 into a separate partition. As we have already seen, this cannot be done using `ALTER TABLE ... ADD PARTITION`. However, you can use another partition-related extension to `ALTER TABLE` to accomplish this:

```
ALTER TABLE members REORGANIZE PARTITION n0 INTO (
```

```
    PARTITION s0 VALUES LESS THAN (1960),
    PARTITION s1 VALUES LESS THAN (1970)
);
```

In effect, this command splits partition `p0` into two new partitions `s0` and `s1`. It also moves the data that was stored in `p0` into the new partitions according to the rules embodied in the two `PARTITION ... VALUES ...` clauses, so that `s0` contains only those records for which `YEAR(dob)` is less than 1960 and `s1` contains those rows in which `YEAR(dob)` is greater than or equal to 1960 but less than 1970.

A `REORGANIZE PARTITION` clause may also be used for merging adjacent partitions. You can reverse the effect of the previous statement on the `members` table as shown here:

```
ALTER TABLE members REORGANIZE PARTITION s0,s1 INTO (
    PARTITION p0 VALUES LESS THAN (1970)
);
```

No data is lost in splitting or merging partitions using `REORGANIZE PARTITION`. In executing the above statement, MySQL moves all of the records that were stored in partitions `s0` and `s1` into partition `p0`.

The general syntax for `REORGANIZE PARTITION` is shown here:

```
ALTER TABLE tbl_name
    REORGANIZE PARTITION partition_list
    INTO (partition_definitions);
```

Here, *`tbl_name`* is the name of the partitioned table, and *`partition_list`* is a comma-separated list of names of one or more existing partitions to be changed. *`partition_definitions`* is a comma-separated list of new partition definitions, which follow the same rules as for the *`partition_definitions`* list used in `CREATE TABLE`. You are not limited to merging several partitions into one, or to splitting one partition into many, when using `REORGANIZE PARTITION`. For example, you can reorganize all four partitions of the `members` table into two, like this:

```
ALTER TABLE members REORGANIZE PARTITION p0,p1,p2,p3 INTO (
    PARTITION m0 VALUES LESS THAN (1980),
    PARTITION m1 VALUES LESS THAN (2000)
);
```

You can also use `REORGANIZE PARTITION` with tables that are partitioned by `LIST`. Let us return to the problem of adding a new partition to the list-partitioned `tt` table and failing because the new partition had a value that was already present in the value-list of one of the existing partitions. We can handle this by adding a partition that contains only nonconflicting values, and then reorganizing the new partition and the existing one so that the value which was stored in the existing one is now moved to the new one:

```
ALTER TABLE tt ADD PARTITION (PARTITION np VALUES IN (4, 8));
ALTER TABLE tt REORGANIZE PARTITION p1,np INTO (
    PARTITION p1 VALUES IN (6, 18),
    PARTITION np VALUES in (4, 8, 12)
);
```

Here are some key points to keep in mind when using `ALTER TABLE ... REORGANIZE PARTITION` to repartition tables that are partitioned by `RANGE` or `LIST`:

- The `PARTITION` options used to determine the new partitioning scheme are subject to the same rules as those used with a `CREATE TABLE` statement.

  A new `RANGE` partitioning scheme cannot have any overlapping ranges; a new `LIST` partitioning scheme cannot have any overlapping sets of values.

- The combination of partitions in the *`partition_definitions`* list should account for the same range or set of values overall as the combined partitions named in the *`partition_list`*.

For example, partitions `p1` and `p2` together cover the years 1980 through 1999 in the `members` table used as an example in this section. Any reorganization of these two partitions should cover the same range of years overall.

- For tables partitioned by `RANGE`, you can reorganize only adjacent partitions; you cannot skip over range partitions.

  For instance, you could not reorganize the example `members` table using a statement beginning with `ALTER TABLE members REORGANIZE PARTITION p0,p2 INTO ...` because `p0` covers the years prior to 1970 and `p2` the years from 1990 through 1999 inclusive, so these are not adjacent partitions. (You cannot skip partition `p1` in this case.)

- You cannot use `REORGANIZE PARTITION` to change the type of partitioning used by the table; for example, you cannot change `RANGE` partitions to `HASH` partitions or the reverse. You also cannot use this statement to change the partitioning expression or column. To accomplish either of these tasks without dropping and re-creating the table, you can use `ALTER TABLE ... PARTITION BY ...`, as shown here:

```
ALTER TABLE members
    PARTITION BY HASH( YEAR(dob) )
    PARTITIONS 8;
```

22.3.2 Management of HASH and KEY Partitions

Tables which are partitioned by hash or by key are very similar to one another with regard to making changes in a partitioning setup, and both differ in a number of ways from tables which have been partitioned by range or list. For that reason, this section addresses the modification of tables partitioned by hash or by key only. For a discussion of adding and dropping of partitions of tables that are partitioned by range or list, see Section 22.3.1, “Management of RANGE and LIST Partitions”.

You cannot drop partitions from tables that are partitioned by `HASH` or `KEY` in the same way that you can from tables that are partitioned by `RANGE` or `LIST`. However, you can merge `HASH` or `KEY` partitions using `ALTER TABLE ... COALESCE PARTITION`. Suppose that a `clients` table containing data about clients is divided into 12 partitions, created as shown here:

```
CREATE TABLE clients (
    id INT,
    fname VARCHAR(30),
    lname VARCHAR(30),
    signed DATE
)
PARTITION BY HASH( MONTH(signed) )
PARTITIONS 12;
```

To reduce the number of partitions from 12 to 8, execute the following `ALTER TABLE` statement:

```
mysql> ALTER TABLE clients COALESCE PARTITION 4;
Query OK, 0 rows affected (0.02 sec)
```

`COALESCE` works equally well with tables that are partitioned by `HASH`, `KEY`, `LINEAR HASH`, or `LINEAR KEY`. Here is an example similar to the previous one, differing only in that the table is partitioned by `LINEAR KEY`:

```
mysql> CREATE TABLE clients_lk (
           id INT,
           fname VARCHAR(30),
           lname VARCHAR(30),
           signed DATE
       )
       PARTITION BY LINEAR KEY(signed)
```

```
      PARTITIONS 12;
Query OK, 0 rows affected (0.03 sec)

mysql> ALTER TABLE clients_lk COALESCE PARTITION 4;
Query OK, 0 rows affected (0.06 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

The number following `COALESCE PARTITION` is the number of partitions to merge into the remainder—in other words, it is the number of partitions to remove from the table.

Attempting to remove more partitions than are in the table results in an error like this one:

```
mysql> ALTER TABLE clients COALESCE PARTITION 18;
ERROR 1478 (HY000): Cannot remove all partitions, use DROP TABLE instead
```

To increase the number of partitions for the `clients` table from 12 to 18, use `ALTER TABLE ... ADD PARTITION` as shown here:

```
ALTER TABLE clients ADD PARTITION PARTITIONS 6;
```

22.3.3 Exchanging Partitions and Subpartitions with Tables

In MySQL 8.0, it is possible to exchange a table partition or subpartition with a table using `ALTER TABLE pt EXCHANGE PARTITION p WITH TABLE nt`, where *`pt`* is the partitioned table and *`p`* is the partition or subpartition of *`pt`* to be exchanged with unpartitioned table *`nt`*, provided that the following statements are true:

1. Table *`nt`* is not itself partitioned.
2. Table *`nt`* is not a temporary table.
3. The structures of tables *`pt`* and *`nt`* are otherwise identical.
4. Table `nt` contains no foreign key references, and no other table has any foreign keys that refer to `nt`.
5. There are no rows in *`nt`* that lie outside the boundaries of the partition definition for *`p`*. This condition does not apply if `WITHOUT VALIDATION` is used.

In addition to the `ALTER`, `INSERT`, and `CREATE` privileges usually required for `ALTER TABLE` statements, you must have the `DROP` privilege to perform `ALTER TABLE ... EXCHANGE PARTITION`.

You should also be aware of the following effects of `ALTER TABLE ... EXCHANGE PARTITION`:

- Executing `ALTER TABLE ... EXCHANGE PARTITION` does not invoke any triggers on either the partitioned table or the table to be exchanged.
- Any `AUTO_INCREMENT` columns in the exchanged table are reset.
- The `IGNORE` keyword has no effect when used with `ALTER TABLE ... EXCHANGE PARTITION`.

The syntax for `ALTER TABLE ... EXCHANGE PARTITION` is shown here, where *`pt`* is the partitioned table, *`p`* is the partition (or subpartition) to be exchanged, and *`nt`* is the nonpartitioned table to be exchanged with *`p`*:

```
ALTER TABLE pt
    EXCHANGE PARTITION p
    WITH TABLE nt;
```

Optionally, you can append `WITH VALIDATION` or `WITHOUT VALIDATION`. When `WITHOUT VALIDATION` is specified, the `ALTER TABLE ... EXCHANGE PARTITION` operation does not

perform any row-by-row validation when exchanging a partition a nonpartitioned table, allowing database administrators to assume responsibility for ensuring that rows are within the boundaries of the partition definition. `WITH VALIDATION` is the default.

One and only one partition or subpartition may be exchanged with one and only one nonpartitioned table in a single `ALTER TABLE EXCHANGE PARTITION` statement. To exchange multiple partitions or subpartitions, use multiple `ALTER TABLE EXCHANGE PARTITION` statements. `EXCHANGE PARTITION` may not be combined with other `ALTER TABLE` options. The partitioning and (if applicable) subpartitioning used by the partitioned table may be of any type or types supported in MySQL 8.0.

Exchanging a Partition with a Nonpartitioned Table

Suppose that a partitioned table `e` has been created and populated using the following SQL statements:

```
CREATE TABLE e (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30)
)
    PARTITION BY RANGE (id) (
        PARTITION p0 VALUES LESS THAN (50),
        PARTITION p1 VALUES LESS THAN (100),
        PARTITION p2 VALUES LESS THAN (150),
        PARTITION p3 VALUES LESS THAN (MAXVALUE)
);

INSERT INTO e VALUES
    (1669, "Jim", "Smith"),
    (337, "Mary", "Jones"),
    (16, "Frank", "White"),
    (2005, "Linda", "Black");
```

Now we create a nonpartitioned copy of `e` named `e2`. This can be done using the `mysql` client as shown here:

```
mysql> CREATE TABLE e2 LIKE e;
Query OK, 0 rows affected (0.04 sec)

mysql> ALTER TABLE e2 REMOVE PARTITIONING;
Query OK, 0 rows affected (0.07 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

You can see which partitions in table `e` contain rows by querying the `INFORMATION_SCHEMA.PARTITIONS` table, like this:

```
mysql> SELECT PARTITION_NAME, TABLE_ROWS
           FROM INFORMATION_SCHEMA.PARTITIONS
           WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          1 |
| p1             |          0 |
| p2             |          0 |
| p3             |          3 |
+----------------+------------+
2 rows in set (0.00 sec)
```

Note

For partitioned `InnoDB` tables, the row count given in the `TABLE_ROWS` column of the `INFORMATION_SCHEMA.PARTITIONS` table is only an estimated value used in SQL optimization, and is not always exact.

To exchange partition `p0` in table `e` with table `e2`, you can use `ALTER TABLE`, as shown here:

```
mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e2;
Query OK, 0 rows affected (0.04 sec)
```

More precisely, the statement just issued causes any rows found in the partition to be swapped with those found in the table. You can observe how this has happened by querying the `INFORMATION_SCHEMA.PARTITIONS` table, as before. The table row that was previously found in partition `p0` is no longer present:

```
mysql> SELECT PARTITION_NAME, TABLE_ROWS
           FROM INFORMATION_SCHEMA.PARTITIONS
           WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          0 |
| p1             |          0 |
| p2             |          0 |
| p3             |          3 |
+----------------+------------+
4 rows in set (0.00 sec)
```

If you query table `e2`, you can see that the “missing” row can now be found there:

```
mysql> SELECT * FROM e2;
+----+-------+-------+
| id | fname | lname |
+----+-------+-------+
| 16 | Frank | White |
+----+-------+-------+
1 row in set (0.00 sec)
```

The table to be exchanged with the partition does not necessarily have to be empty. To demonstrate this, we first insert a new row into table `e`, making sure that this row is stored in partition `p0` by choosing an `id` column value that is less than 50, and verifying this afterward by querying the `PARTITIONS` table:

```
mysql> INSERT INTO e VALUES (41, "Michael", "Green");
Query OK, 1 row affected (0.05 sec)

mysql> SELECT PARTITION_NAME, TABLE_ROWS
           FROM INFORMATION_SCHEMA.PARTITIONS
           WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          1 |
| p1             |          0 |
| p2             |          0 |
| p3             |          3 |
+----------------+------------+
4 rows in set (0.00 sec)
```

Now we once again exchange partition `p0` with table `e2` using the same `ALTER TABLE` statement as previously:

```
mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e2;
Query OK, 0 rows affected (0.28 sec)
```

The output of the following queries shows that the table row that was stored in partition `p0` and the table row that was stored in table `e2`, prior to issuing the `ALTER TABLE` statement, have now switched places:

```
mysql> SELECT * FROM e;
+------+-------+-------+
| id   | fname | lname |
+------+-------+-------+
|   16 | Frank | White |
| 1669 | Jim   | Smith |
|  337 | Mary  | Jones |
| 2005 | Linda | Black |
+------+-------+-------+
4 rows in set (0.00 sec)

mysql> SELECT PARTITION_NAME, TABLE_ROWS
           FROM INFORMATION_SCHEMA.PARTITIONS
           WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          1 |
| p1             |          0 |
| p2             |          0 |
| p3             |          3 |
+----------------+------------+
4 rows in set (0.00 sec)

mysql> SELECT * FROM e2;
+----+---------+-------+
| id | fname   | lname |
+----+---------+-------+
| 41 | Michael | Green |
+----+---------+-------+
1 row in set (0.00 sec)
```

Nonmatching Rows

You should keep in mind that any rows found in the nonpartitioned table prior to issuing the `ALTER TABLE ... EXCHANGE PARTITION` statement must meet the conditions required for them to be stored in the target partition; otherwise, the statement fails. To see how this occurs, first insert a row into `e2` that is outside the boundaries of the partition definition for partition `p0` of table `e`. For example, insert a row with an `id` column value that is too large; then, try to exchange the table with the partition again:

```
mysql> INSERT INTO e2 VALUES (51, "Ellen", "McDonald");
Query OK, 1 row affected (0.08 sec)

mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e2;
ERROR 1707 (HY000): Found row that does not match the partition
```

Only the `WITHOUT VALIDATION` option would permit this operation to succeed:

```
mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e2 WITHOUT VALIDATION;
Query OK, 0 rows affected (0.02 sec)
```

When a partition is exchanged with a table that contains rows that do not match the partition definition, it is the responsibility of the database administrator to fix the non-matching rows, which can be performed using `REPAIR TABLE` or `ALTER TABLE ... REPAIR PARTITION`.

Exchanging Partitions Without Row-By-Row Validation

To avoid time consuming validation when exchanging a partition with a table that has many rows, it is possible to skip the row-by-row validation step by appending `WITHOUT VALIDATION` to the `ALTER TABLE ... EXCHANGE PARTITION` statement.

The following example compares the difference between execution times when exchanging a partition with a nonpartitioned table, with and without validation. The partitioned table (table `e`) contains two

partitions of 1 million rows each. The rows in p0 of table e are removed and p0 is exchanged with a nonpartitioned table of 1 million rows. The `WITH VALIDATION` operation takes 0.74 seconds. By comparison, the `WITHOUT VALIDATION` operation takes 0.01 seconds.

```
# Create a partitioned table with 1 million rows in each partition

CREATE TABLE e (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30)
)
    PARTITION BY RANGE (id) (
        PARTITION p0 VALUES LESS THAN (1000001),
        PARTITION p1 VALUES LESS THAN (2000001),
);

mysql> SELECT COUNT(*) FROM e;
| COUNT(*) |
+----------+
|  2000000 |
+----------+
1 row in set (0.27 sec)

# View the rows in each partition

SELECT PARTITION_NAME, TABLE_ROWS FROM INFORMATION_SCHEMA.PARTITIONS WHERE TABLE_NAME = 'e';
+----------------+-------------+
| PARTITION_NAME | TABLE_ROWS  |
+----------------+-------------+
| p0             |     1000000 |
| p1             |     1000000 |
+----------------+-------------+
2 rows in set (0.00 sec)

# Create a nonpartitioned table of the same structure and populate it with 1 million rows

CREATE TABLE e2 (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30)
);

mysql> SELECT COUNT(*) FROM e2;
+----------+
| COUNT(*) |
+----------+
|  1000000 |
+----------+
1 row in set (0.24 sec)

# Create another nonpartitioned table of the same structure and populate it with 1 million rows

CREATE TABLE e3 (
    id INT NOT NULL,
    fname VARCHAR(30),
    lname VARCHAR(30)
);

mysql> SELECT COUNT(*) FROM e3;
+----------+
| COUNT(*) |
+----------+
|  1000000 |
+----------+
1 row in set (0.25 sec)

# Drop the rows from p0 of table e

mysql> DELETE FROM e WHERE id < 1000001;
Query OK, 1000000 rows affected (5.55 sec)
```

```
# Confirm that there are no rows in partition p0

mysql> SELECT PARTITION_NAME, TABLE_ROWS FROM INFORMATION_SCHEMA.PARTITIONS WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          0 |
| p1             |    1000000 |
+----------------+------------+
2 rows in set (0.00 sec)

# Exchange partition p0 of table e with the table e2 'WITH VALIDATION'

mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e2 WITH VALIDATION;
Query OK, 0 rows affected (0.74 sec)

# Confirm that the partition was exchanged with table e2

mysql> SELECT PARTITION_NAME, TABLE_ROWS FROM INFORMATION_SCHEMA.PARTITIONS WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |    1000000 |
| p1             |    1000000 |
+----------------+------------+
2 rows in set (0.00 sec)

# Once again, drop the rows from p0 of table e

mysql> DELETE FROM e WHERE id < 1000001;
Query OK, 1000000 rows affected (5.55 sec)

# Confirm that there are no rows in partition p0

mysql> SELECT PARTITION_NAME, TABLE_ROWS FROM INFORMATION_SCHEMA.PARTITIONS WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |          0 |
| p1             |    1000000 |
+----------------+------------+
2 rows in set (0.00 sec)

# Exchange partition p0 of table e with the table e3 'WITHOUT VALIDATION'

mysql> ALTER TABLE e EXCHANGE PARTITION p0 WITH TABLE e3 WITHOUT VALIDATION;
Query OK, 0 rows affected (0.01 sec)

# Confirm that the partition was exchanged with table e3

mysql> SELECT PARTITION_NAME, TABLE_ROWS FROM INFORMATION_SCHEMA.PARTITIONS WHERE TABLE_NAME = 'e';
+----------------+------------+
| PARTITION_NAME | TABLE_ROWS |
+----------------+------------+
| p0             |    1000000 |
| p1             |    1000000 |
+----------------+------------+
2 rows in set (0.00 sec)
```

If a partition is exchanged with a table that contains rows that do not match the partition definition, it is the responsibility of the database administrator to fix the non-matching rows, which can be performed using `REPAIR TABLE` or `ALTER TABLE ... REPAIR PARTITION`.

Exchanging a Subpartition with a Nonpartitioned Table

You can also exchange a subpartition of a subpartitioned table (see Section 22.2.6, “Subpartitioning”) with a nonpartitioned table using an `ALTER TABLE ... EXCHANGE PARTITION` statement. In the following example, we first create a table `es` that is partitioned by `RANGE` and subpartitioned by `KEY`,

populate this table as we did table `e`, and then create an empty, nonpartitioned copy `es2` of the table, as shown here:

```
mysql> CREATE TABLE es (
    ->     id INT NOT NULL,
    ->     fname VARCHAR(30),
    ->     lname VARCHAR(30)
    -> )
    ->     PARTITION BY RANGE (id)
    ->     SUBPARTITION BY KEY (lname)
    ->     SUBPARTITIONS 2 (
    ->         PARTITION p0 VALUES LESS THAN (50),
    ->         PARTITION p1 VALUES LESS THAN (100),
    ->         PARTITION p2 VALUES LESS THAN (150),
    ->         PARTITION p3 VALUES LESS THAN (MAXVALUE)
    ->     );
Query OK, 0 rows affected (2.76 sec)

mysql> INSERT INTO es VALUES
    ->     (1669, "Jim", "Smith"),
    ->     (337, "Mary", "Jones"),
    ->     (16, "Frank", "White"),
    ->     (2005, "Linda", "Black");
Query OK, 4 rows affected (0.04 sec)
Records: 4  Duplicates: 0  Warnings: 0

mysql> CREATE TABLE es2 LIKE es;
Query OK, 0 rows affected (1.27 sec)

mysql> ALTER TABLE es2 REMOVE PARTITIONING;
Query OK, 0 rows affected (0.70 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

Although we did not explicitly name any of the subpartitions when creating table `es`, we can obtain generated names for these by including the `SUBPARTITION_NAME` column of the `PARTITIONS` table from `INFORMATION_SCHEMA` when selecting from that table, as shown here:

```
mysql> SELECT PARTITION_NAME, SUBPARTITION_NAME, TABLE_ROWS
    ->     FROM INFORMATION_SCHEMA.PARTITIONS
    ->     WHERE TABLE_NAME = 'es';
+----------------+-------------------+------------+
| PARTITION_NAME | SUBPARTITION_NAME | TABLE_ROWS |
+----------------+-------------------+------------+
| p0             | p0sp0             |          1 |
| p0             | p0sp1             |          0 |
| p1             | p1sp0             |          0 |
| p1             | p1sp1             |          0 |
| p2             | p2sp0             |          0 |
| p2             | p2sp1             |          0 |
| p3             | p3sp0             |          3 |
| p3             | p3sp1             |          0 |
+----------------+-------------------+------------+
8 rows in set (0.00 sec)
```

The following `ALTER TABLE` statement exchanges subpartition `p3sp0` in table `es` with the nonpartitioned table `es2`:

```
mysql> ALTER TABLE es EXCHANGE PARTITION p3sp0 WITH TABLE es2;
Query OK, 0 rows affected (0.29 sec)
```

You can verify that the rows were exchanged by issuing the following queries:

```
mysql> SELECT PARTITION_NAME, SUBPARTITION_NAME, TABLE_ROWS
    ->     FROM INFORMATION_SCHEMA.PARTITIONS
    ->     WHERE TABLE_NAME = 'es';
+----------------+-------------------+------------+
```

```
| PARTITION_NAME | SUBPARTITION_NAME | TABLE_ROWS |
+----------------+-------------------+------------+
| p0             | p0sp0             |          1 |
| p0             | p0sp1             |          0 |
| p1             | p1sp0             |          0 |
| p1             | p1sp1             |          0 |
| p2             | p2sp0             |          0 |
| p2             | p2sp1             |          0 |
| p3             | p3sp0             |          0 |
| p3             | p3sp1             |          0 |
+----------------+-------------------+------------+
8 rows in set (0.00 sec)

mysql> SELECT * FROM es2;
+------+-------+-------+
| id   | fname | lname |
+------+-------+-------+
| 1669 | Jim   | Smith |
|  337 | Mary  | Jones |
| 2005 | Linda | Black |
+------+-------+-------+
3 rows in set (0.00 sec)
```

If a table is subpartitioned, you can exchange only a subpartition of the table—not an entire partition—with an unpartitioned table, as shown here:

```
mysql> ALTER TABLE es EXCHANGE PARTITION p3 WITH TABLE es2;
ERROR 1704 (HY000): Subpartitioned table, use subpartition instead of partition
```

Table structures are compared in a strict fashion; the number, order, names, and types of columns and indexes of the partitioned table and the nonpartitioned table must match exactly. In addition, both tables must use the same storage engine:

```
mysql> CREATE TABLE es3 LIKE e;
Query OK, 0 rows affected (1.31 sec)

mysql> ALTER TABLE es3 REMOVE PARTITIONING;
Query OK, 0 rows affected (0.53 sec)
Records: 0  Duplicates: 0  Warnings: 0

mysql> SHOW CREATE TABLE es3\G
*************************** 1. row ***************************
       Table: es3
Create Table: CREATE TABLE `es3` (
  `id` int(11) NOT NULL,
  `fname` varchar(30) DEFAULT NULL,
  `lname` varchar(30) DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4
1 row in set (0.00 sec)

mysql> ALTER TABLE es3 ENGINE = MyISAM;
Query OK, 0 rows affected (0.15 sec)
Records: 0  Duplicates: 0  Warnings: 0

mysql> ALTER TABLE es EXCHANGE PARTITION p3sp0 WITH TABLE es3;
ERROR 1497 (HY000): The mix of handlers in the partitions is not allowed in this version of MySQL
```

22.3.4 Maintenance of Partitions

A number of table and partition maintenance tasks can be carried out on partitioned tables using SQL statements intended for such purposes.

Table maintenance of partitioned tables can be accomplished using the statements `CHECK TABLE`, `OPTIMIZE TABLE`, `ANALYZE TABLE`, and `REPAIR TABLE`, which are supported for partitioned tables.

You can use a number of extensions to `ALTER TABLE` for performing operations of this type on one or more partitions directly, as described in the following list:

- **Rebuilding partitions.** Rebuilds the partition; this has the same effect as dropping all records stored in the partition, then reinserting them. This can be useful for purposes of defragmentation.

  Example:

  ```
  ALTER TABLE t1 REBUILD PARTITION p0, p1;
  ```

- **Optimizing partitions.** If you have deleted a large number of rows from a partition or if you have made many changes to a partitioned table with variable-length rows (that is, having `VARCHAR`, `BLOB`, or `TEXT` columns), you can use `ALTER TABLE ... OPTIMIZE PARTITION` to reclaim any unused space and to defragment the partition data file.

  Example:

  ```
  ALTER TABLE t1 OPTIMIZE PARTITION p0, p1;
  ```

  Using `OPTIMIZE PARTITION` on a given partition is equivalent to running `CHECK PARTITION`, `ANALYZE PARTITION`, and `REPAIR PARTITION` on that partition.

  Some MySQL storage engines, including `InnoDB`, do not support per-partition optimization; in these cases, `ALTER TABLE ... OPTIMIZE PARTITION` analyzes and rebuilds the entire table, and causes an appropriate warning to be issued. (Bug #11751825, Bug #42822) Use `ALTER TABLE ... REBUILD PARTITION` and `ALTER TABLE ... ANALYZE PARTITION` instead, to avoid this issue.

- **Analyzing partitions.** This reads and stores the key distributions for partitions.

  Example:

  ```
  ALTER TABLE t1 ANALYZE PARTITION p3;
  ```

- **Repairing partitions.** This repairs corrupted partitions.

  Example:

  ```
  ALTER TABLE t1 REPAIR PARTITION p0,p1;
  ```

  Normally, `REPAIR PARTITION` fails when the partition contains duplicate key errors. You can use `ALTER IGNORE TABLE` with this option, in which case all rows that cannot be moved due to the presence of duplicate keys are removed from the partition (Bug #16900947).

- **Checking partitions.** You can check partitions for errors in much the same way that you can use `CHECK TABLE` with nonpartitioned tables.

  Example:

  ```
  ALTER TABLE trb3 CHECK PARTITION p1;
  ```

  This command will tell you whether the data or indexes in partition `p1` of table `t1` are corrupted. If this is the case, use `ALTER TABLE ... REPAIR PARTITION` to repair the partition.

  Normally, `CHECK PARTITION` fails when the partition contains duplicate key errors. You can use `ALTER IGNORE TABLE` with this option, in which case the statement returns the contents of each row in the partition where a duplicate key violation is found. Only the values for the columns in the partitioning expression for the table are reported. (Bug #16900947)

Each of the statements in the list just shown also supports the keyword `ALL` in place of the list of partition names. Using `ALL` causes the statement to act on all partitions in the table.

You can also truncate partitions using `ALTER TABLE ... TRUNCATE PARTITION`. This statement can be used to delete all rows from one or more partitions in much the same way that `TRUNCATE TABLE` deletes all rows from a table.

`ALTER TABLE ... TRUNCATE PARTITION ALL` truncates all partitions in the table.

22.3.5 Obtaining Information About Partitions

This section discusses obtaining information about existing partitions, which can be done in a number of ways. Methods of obtaining such information include the following:

- Using the `SHOW CREATE TABLE` statement to view the partitioning clauses used in creating a partitioned table.
- Using the `SHOW TABLE STATUS` statement to determine whether a table is partitioned.
- Querying the `INFORMATION_SCHEMA.PARTITIONS` table.
- Using the statement `EXPLAIN SELECT` to see which partitions are used by a given `SELECT`.

As discussed elsewhere in this chapter, `SHOW CREATE TABLE` includes in its output the `PARTITION BY` clause used to create a partitioned table. For example:

```
mysql> SHOW CREATE TABLE trb3\G
*************************** 1. row ***************************
       Table: trb3
Create Table: CREATE TABLE `trb3` (
  `id` int(11) DEFAULT NULL,
  `name` varchar(50) DEFAULT NULL,
  `purchased` date DEFAULT NULL
) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4
/*!50100 PARTITION BY RANGE (YEAR(purchased))
(PARTITION p0 VALUES LESS THAN (1990) ENGINE = InnoDB,
 PARTITION p1 VALUES LESS THAN (1995) ENGINE = InnoDB,
 PARTITION p2 VALUES LESS THAN (2000) ENGINE = InnoDB,
 PARTITION p3 VALUES LESS THAN (2005) ENGINE = InnoDB) */
0 row in set (0.00 sec)
```

The output from `SHOW TABLE STATUS` for partitioned tables is the same as that for nonpartitioned tables, except that the `Create_options` column contains the string `partitioned`. The `Engine` column contains the name of the storage engine used by all partitions of the table. (See Section 13.7.5.36, “SHOW TABLE STATUS Syntax”, for more information about this statement.)

You can also obtain information about partitions from `INFORMATION_SCHEMA`, which contains a `PARTITIONS` table. See Section 24.13, “The INFORMATION_SCHEMA PARTITIONS Table”.

It is possible to determine which partitions of a partitioned table are involved in a given `SELECT` query using `EXPLAIN`. The `partitions` column in the `EXPLAIN` output lists the partitions from which records would be matched by the query.

Suppose that a table `trb1` is created and populated as follows:

```
CREATE TABLE trb1 (id INT, name VARCHAR(50), purchased DATE)
    PARTITION BY RANGE(id)
    (
        PARTITION p0 VALUES LESS THAN (3),
        PARTITION p1 VALUES LESS THAN (7),
        PARTITION p2 VALUES LESS THAN (9),
        PARTITION p3 VALUES LESS THAN (11)
    );

INSERT INTO trb1 VALUES
    (1, 'desk organiser', '2003-10-15'),
    (2, 'CD player', '1993-11-05'),
```

```
    (3, 'TV set', '1996-03-10'),
    (4, 'bookcase', '1982-01-10'),
    (5, 'exercise bike', '2004-05-09'),
    (6, 'sofa', '1987-06-05'),
    (7, 'popcorn maker', '2001-11-22'),
    (8, 'aquarium', '1992-08-04'),
    (9, 'study desk', '1984-09-16'),
    (10, 'lava lamp', '1998-12-25');
```

You can see which partitions are used in a query such as `SELECT * FROM trb1;`, as shown here:

```
mysql> EXPLAIN SELECT * FROM trb1\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: trb1
   partitions: p0,p1,p2,p3
         type: ALL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 10
        Extra: Using filesort
```

In this case, all four partitions are searched. However, when a limiting condition making use of the partitioning key is added to the query, you can see that only those partitions containing matching values are scanned, as shown here:

```
mysql> EXPLAIN SELECT * FROM trb1 WHERE id < 5\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: trb1
   partitions: p0,p1
         type: ALL
possible_keys: NULL
          key: NULL
      key_len: NULL
          ref: NULL
         rows: 10
        Extra: Using where
```

`EXPLAIN` also provides information about keys used and possible keys:

```
mysql> ALTER TABLE trb1 ADD PRIMARY KEY (id);
Query OK, 10 rows affected (0.03 sec)
Records: 10  Duplicates: 0  Warnings: 0

mysql> EXPLAIN SELECT * FROM trb1 WHERE id < 5\G
*************************** 1. row ***************************
           id: 1
  select_type: SIMPLE
        table: trb1
   partitions: p0,p1
         type: range
possible_keys: PRIMARY
          key: PRIMARY
      key_len: 4
          ref: NULL
         rows: 7
        Extra: Using where
```

If `EXPLAIN` is used to examine a query against a nonpartitioned table, no error is produced, but the value of the `partitions` column is always `NULL`.

The `rows` column of `EXPLAIN` output displays the total number of rows in the table.

See also Section 13.8.2, “EXPLAIN Syntax”.

22.4 Partition Pruning

The optimization known as *partition pruning* is based on a relatively simple concept which can be described as “Do not scan partitions where there can be no matching values”. Suppose a partitioned table `t1` is created by this statement:

```
CREATE TABLE t1 (
    fname VARCHAR(50) NOT NULL,
    lname VARCHAR(50) NOT NULL,
    region_code TINYINT UNSIGNED NOT NULL,
    dob DATE NOT NULL
)
PARTITION BY RANGE( region_code ) (
    PARTITION p0 VALUES LESS THAN (64),
    PARTITION p1 VALUES LESS THAN (128),
    PARTITION p2 VALUES LESS THAN (192),
    PARTITION p3 VALUES LESS THAN MAXVALUE
);
```

Suppose that you wish to obtain results from a `SELECT` statement such as this one:

```
SELECT fname, lname, region_code, dob
    FROM t1
    WHERE region_code > 125 AND region_code < 130;
```

It is easy to see that none of the rows which ought to be returned are in either of the partitions `p0` or `p3`; that is, we need search only in partitions `p1` and `p2` to find matching rows. By limiting the search, it is possible to expend much less time and effort in finding matching rows than by scanning all partitions in the table. This “cutting away” of unneeded partitions is known as *pruning*. When the optimizer can make use of partition pruning in performing this query, execution of the query can be an order of magnitude faster than the same query against a nonpartitioned table containing the same column definitions and data.

The optimizer can perform pruning whenever a `WHERE` condition can be reduced to either one of the following two cases:

- *`partition_column = constant`*
- *`partition_column`* `IN (`*`constant1, constant2, ..., constantN`*`)`

In the first case, the optimizer simply evaluates the partitioning expression for the value given, determines which partition contains that value, and scans only this partition. In many cases, the equal sign can be replaced with another arithmetic comparison, including `<`, `>`, `<=`, `>=`, and `<>`. Some queries using `BETWEEN` in the `WHERE` clause can also take advantage of partition pruning. See the examples later in this section.

In the second case, the optimizer evaluates the partitioning expression for each value in the list, creates a list of matching partitions, and then scans only the partitions in this partition list.

`SELECT`, `DELETE`, and `UPDATE` statements support partition pruning. `INSERT` statements currently cannot be pruned.

Pruning can also be applied to short ranges, which the optimizer can convert into equivalent lists of values. For instance, in the previous example, the `WHERE` clause can be converted to `WHERE region_code IN (126, 127, 128, 129)`. Then the optimizer can determine that the first two values in the list are found in partition `p1`, the remaining two values in partition `p2`, and that the other partitions contain no relevant values and so do not need to be searched for matching rows.

The optimizer can also perform pruning for `WHERE` conditions that involve comparisons of the preceding types on multiple columns for tables that use `RANGE COLUMNS` or `LIST COLUMNS` partitioning.

This type of optimization can be applied whenever the partitioning expression consists of an equality or a range which can be reduced to a set of equalities, or when the partitioning expression represents an increasing or decreasing relationship. Pruning can also be applied for tables partitioned on a `DATE` or `DATETIME` column when the partitioning expression uses the `YEAR()` or `TO_DAYS()` function. Pruning can also be applied for such tables when the partitioning expression uses the `TO_SECONDS()` function.

Suppose that table `t2`, partitioned on a `DATE` column, is created using the statement shown here:

```
CREATE TABLE t2 (
    fname VARCHAR(50) NOT NULL,
    lname VARCHAR(50) NOT NULL,
    region_code TINYINT UNSIGNED NOT NULL,
    dob DATE NOT NULL
)
PARTITION BY RANGE( YEAR(dob) ) (
    PARTITION d0 VALUES LESS THAN (1970),
    PARTITION d1 VALUES LESS THAN (1975),
    PARTITION d2 VALUES LESS THAN (1980),
    PARTITION d3 VALUES LESS THAN (1985),
    PARTITION d4 VALUES LESS THAN (1990),
    PARTITION d5 VALUES LESS THAN (2000),
    PARTITION d6 VALUES LESS THAN (2005),
    PARTITION d7 VALUES LESS THAN MAXVALUE
);
```

The following statements using `t2` can make of use partition pruning:

```
SELECT * FROM t2 WHERE dob = '1982-06-23';

UPDATE t2 SET region_code = 8 WHERE dob BETWEEN '1991-02-15' AND '1997-04-25';

DELETE FROM t2 WHERE dob >= '1984-06-21' AND dob <= '1999-06-21'
```

In the case of the last statement, the optimizer can also act as follows:

1. *Find the partition containing the low end of the range*.

   `YEAR('1984-06-21')` yields the value `1984`, which is found in partition `d3`.

2. *Find the partition containing the high end of the range*.

   `YEAR('1999-06-21')` evaluates to `1999`, which is found in partition `d5`.

3. *Scan only these two partitions and any partitions that may lie between them*.

   In this case, this means that only partitions `d3`, `d4`, and `d5` are scanned. The remaining partitions may be safely ignored (and are ignored).

Important

Invalid `DATE` and `DATETIME` values referenced in the `WHERE` condition of a statement against a partitioned table are treated as `NULL`. This means that a query such as `SELECT * FROM partitioned_table WHERE date_column < '2008-12-00'` does not return any values (see Bug #40972).

So far, we have looked only at examples using `RANGE` partitioning, but pruning can be applied with other partitioning types as well.

Consider a table that is partitioned by `LIST`, where the partitioning expression is increasing or decreasing, such as the table `t3` shown here. (In this example, we assume for the sake of brevity that the `region_code` column is limited to values between 1 and 10 inclusive.)

```
CREATE TABLE t3 (
    fname VARCHAR(50) NOT NULL,
    lname VARCHAR(50) NOT NULL,
    region_code TINYINT UNSIGNED NOT NULL,
    dob DATE NOT NULL
)
PARTITION BY LIST(region_code) (
    PARTITION r0 VALUES IN (1, 3),
    PARTITION r1 VALUES IN (2, 5, 8),
    PARTITION r2 VALUES IN (4, 9),
    PARTITION r3 VALUES IN (6, 7, 10)
);
```

For a statement such as `SELECT * FROM t3 WHERE region_code BETWEEN 1 AND 3`, the optimizer determines in which partitions the values 1, 2, and 3 are found (`r0` and `r1`) and skips the remaining ones (`r2` and `r3`).

For tables that are partitioned by `HASH` or `[LINEAR] KEY`, partition pruning is also possible in cases in which the `WHERE` clause uses a simple `=` relation against a column used in the partitioning expression. Consider a table created like this:

```
CREATE TABLE t4 (
    fname VARCHAR(50) NOT NULL,
    lname VARCHAR(50) NOT NULL,
    region_code TINYINT UNSIGNED NOT NULL,
    dob DATE NOT NULL
)
PARTITION BY KEY(region_code)
PARTITIONS 8;
```

A statement that compares a column value with a constant can be pruned:

```
UPDATE t4 WHERE region_code = 7;
```

Pruning can also be employed for short ranges, because the optimizer can turn such conditions into `IN` relations. For example, using the same table `t4` as defined previously, queries such as these can be pruned:

```
SELECT * FROM t4 WHERE region_code > 2 AND region_code < 6;

SELECT * FROM t4 WHERE region_code BETWEEN 3 AND 5;
```

In both these cases, the `WHERE` clause is transformed by the optimizer into `WHERE region_code IN (3, 4, 5)`.

Important

This optimization is used only if the range size is smaller than the number of partitions. Consider this statement:

```
DELETE FROM t4 WHERE region_code BETWEEN 4 AND 12;
```

The range in the `WHERE` clause covers 9 values (4, 5, 6, 7, 8, 9, 10, 11, 12), but `t4` has only 8 partitions. This means that the `DELETE` cannot be pruned.

When a table is partitioned by `HASH` or `[LINEAR] KEY`, pruning can be used only on integer columns. For example, this statement cannot use pruning because `dob` is a `DATE` column:

```
SELECT * FROM t4 WHERE dob >= '2001-04-14' AND dob <= '2005-10-15';
```

However, if the table stores year values in an `INT` column, then a query having `WHERE year_col >= 2001 AND year_col <= 2005` can be pruned.

Tables using a storage engine that provides automatic partitioning, such as the NDB storage engine used by MySQL Cluster (not currently supported in MySQL 8.0), can be pruned if they are explicitly partitioned.

22.5 Partition Selection

Explicit selection of partitions and subpartitions for rows matching a given WHERE condition is supported. Partition selection is similar to partition pruning, in that only specific partitions are checked for matches, but differs in two key respects:

1. The partitions to be checked are specified by the issuer of the statement, unlike partition pruning, which is automatic.
2. Whereas partition pruning applies only to queries, explicit selection of partitions is supported for both queries and a number of DML statements.

SQL statements supporting explicit partition selection are listed here:

- SELECT
- DELETE
- INSERT
- REPLACE
- UPDATE
- LOAD DATA.
- LOAD XML.

The remainder of this section discusses explicit partition selection as it applies generally to the statements just listed, and provides some examples.

Explicit partition selection is implemented using a PARTITION option. For all supported statements, this option uses the syntax shown here:

```
    PARTITION (partition_names)

    partition_names:
        partition_name, ...
```

This option always follows the name of the table to which the partition or partitions belong. *partition_names* is a comma-separated list of partitions or subpartitions to be used. Each name in this list must be the name of an existing partition or subpartition of the specified table; if any of the partitions or subpartitions are not found, the statement fails with an error (partition 'partition_name' doesn't exist). Partitions and subpartitions named in *partition_names* may be listed in any order, and may overlap.

When the PARTITION option is used, only the partitions and subpartitions listed are checked for matching rows. This option can be used in a SELECT statement to determine which rows belong to a given partition. Consider a partitioned table named employees, created and populated using the statements shown here:

```
SET @@SQL_MODE = '';

CREATE TABLE employees  (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    fname VARCHAR(25) NOT NULL,
    lname VARCHAR(25) NOT NULL,
    store_id INT NOT NULL,
    department_id INT NOT NULL
```

```
)
    PARTITION BY RANGE(id)  (
        PARTITION p0 VALUES LESS THAN (5),
        PARTITION p1 VALUES LESS THAN (10),
        PARTITION p2 VALUES LESS THAN (15),
        PARTITION p3 VALUES LESS THAN MAXVALUE
);

INSERT INTO employees VALUES
    ('', 'Bob', 'Taylor', 3, 2), ('', 'Frank', 'Williams', 1, 2),
    ('', 'Ellen', 'Johnson', 3, 4), ('', 'Jim', 'Smith', 2, 4),
    ('', 'Mary', 'Jones', 1, 1), ('', 'Linda', 'Black', 2, 3),
    ('', 'Ed', 'Jones', 2, 1), ('', 'June', 'Wilson', 3, 1),
    ('', 'Andy', 'Smith', 1, 3), ('', 'Lou', 'Waters', 2, 4),
    ('', 'Jill', 'Stone', 1, 4), ('', 'Roger', 'White', 3, 2),
    ('', 'Howard', 'Andrews', 1, 2), ('', 'Fred', 'Goldberg', 3, 3),
    ('', 'Barbara', 'Brown', 2, 3), ('', 'Alice', 'Rogers', 2, 2),
    ('', 'Mark', 'Morgan', 3, 3), ('', 'Karen', 'Cole', 3, 2);
```

You can see which rows are stored in partition `p1` like this:

```
mysql> SELECT * FROM employees PARTITION (p1);
+----+-------+--------+----------+---------------+
| id | fname | lname  | store_id | department_id |
+----+-------+--------+----------+---------------+
|  5 | Mary  | Jones  |        1 |             1 |
|  6 | Linda | Black  |        2 |             3 |
|  7 | Ed    | Jones  |        2 |             1 |
|  8 | June  | Wilson |        3 |             1 |
|  9 | Andy  | Smith  |        1 |             3 |
+----+-------+--------+----------+---------------+
5 rows in set (0.00 sec)
```

The result is the same as obtained by the query `SELECT * FROM employees WHERE id BETWEEN 5 AND 9`.

To obtain rows from multiple partitions, supply their names as a comma-delimited list. For example, `SELECT * FROM employees PARTITION (p1, p2)` returns all rows from partitions `p1` and `p2` while excluding rows from the remaining partitions.

Any valid query against a partitioned table can be rewritten with a `PARTITION` option to restrict the result to one or more desired partitions. You can use `WHERE` conditions, `ORDER BY` and `LIMIT` options, and so on. You can also use aggregate functions with `HAVING` and `GROUP BY` options. Each of the following queries produces a valid result when run on the `employees` table as previously defined:

```
mysql> SELECT * FROM employees PARTITION (p0, p2)
    ->     WHERE lname LIKE 'S%';
+----+-------+-------+----------+---------------+
| id | fname | lname | store_id | department_id |
+----+-------+-------+----------+---------------+
|  4 | Jim   | Smith |        2 |             4 |
| 11 | Jill  | Stone |        1 |             4 |
+----+-------+-------+----------+---------------+
2 rows in set (0.00 sec)

mysql> SELECT id, CONCAT(fname, ' ', lname) AS name
    ->     FROM employees PARTITION (p0) ORDER BY lname;
+----+----------------+
| id | name           |
+----+----------------+
|  3 | Ellen Johnson  |
|  4 | Jim Smith      |
|  1 | Bob Taylor     |
|  2 | Frank Williams |
+----+----------------+
4 rows in set (0.06 sec)

mysql> SELECT store_id, COUNT(department_id) AS c
```

```
    ->     FROM employees PARTITION (p1,p2,p3)
    ->     GROUP BY store_id HAVING c > 4;
+---+----------+
| c | store_id |
+---+----------+
| 5 |        2 |
| 5 |        3 |
+---+----------+
2 rows in set (0.00 sec)
```

Statements using partition selection can be employed with tables using any of the supported partitioning types. When a table is created using `[LINEAR] HASH` or `[LINEAR] KEY` partitioning and the names of the partitions are not specified, MySQL automatically names the partitions `p0`, `p1`, `p2`, ..., `pN-1`, where *`N`* is the number of partitions. For subpartitions not explicitly named, MySQL assigns automatically to the subpartitions in each partition `pX` the names `pXsp0`, `pXsp1`, `pXsp2`, ..., `pXspM-1`, where *`M`* is the number of subpartitions. When executing against this table a `SELECT` (or other SQL statement for which explicit partition selection is allowed), you can use these generated names in a `PARTITION` option, as shown here:

```
mysql> CREATE TABLE employees_sub  (
    ->     id INT NOT NULL AUTO_INCREMENT,
    ->     fname VARCHAR(25) NOT NULL,
    ->     lname VARCHAR(25) NOT NULL,
    ->     store_id INT NOT NULL,
    ->     department_id INT NOT NULL,
    ->     PRIMARY KEY pk (id, lname)
    -> )
    ->     PARTITION BY RANGE(id)
    ->     SUBPARTITION BY KEY (lname)
    ->     SUBPARTITIONS 2 (
    ->         PARTITION p0 VALUES LESS THAN (5),
    ->         PARTITION p1 VALUES LESS THAN (10),
    ->         PARTITION p2 VALUES LESS THAN (15),
    ->         PARTITION p3 VALUES LESS THAN MAXVALUE
    -> );
Query OK, 0 rows affected (1.14 sec)

mysql> INSERT INTO employees_sub   # reuse data in employees table
    ->     SELECT * FROM employees;
Query OK, 18 rows affected (0.09 sec)
Records: 18  Duplicates: 0  Warnings: 0

mysql> SELECT id, CONCAT(fname, ' ', lname) AS name
    ->     FROM employees_sub PARTITION (p2sp1);
+----+---------------+
| id | name          |
+----+---------------+
| 10 | Lou Waters    |
| 14 | Fred Goldberg |
+----+---------------+
2 rows in set (0.00 sec)
```

You may also use a `PARTITION` option in the `SELECT` portion of an `INSERT ... SELECT` statement, as shown here:

```
mysql> CREATE TABLE employees_copy LIKE employees;
Query OK, 0 rows affected (0.28 sec)

mysql> INSERT INTO employees_copy
    ->     SELECT * FROM employees PARTITION (p2);
Query OK, 5 rows affected (0.04 sec)
Records: 5  Duplicates: 0  Warnings: 0

mysql> SELECT * FROM employees_copy;
+----+--------+----------+----------+---------------+
| id | fname  | lname    | store_id | department_id |
+----+--------+----------+----------+---------------+
| 10 | Lou    | Waters   |        2 |             4 |
```

```
| 11 | Jill   | Stone    |         1 |             4 |
| 12 | Roger  | White    |         3 |             2 |
| 13 | Howard | Andrews  |         1 |             2 |
| 14 | Fred   | Goldberg |         3 |             3 |
+----+--------+----------+-----------+---------------+
5 rows in set (0.00 sec)
```

Partition selection can also be used with joins. Suppose we create and populate two tables using the statements shown here:

```
CREATE TABLE stores (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    city VARCHAR(30) NOT NULL
)
    PARTITION BY HASH(id)
    PARTITIONS 2;

INSERT INTO stores VALUES
    ('', 'Nambucca'), ('', 'Uranga'),
    ('', 'Bellingen'), ('', 'Grafton');

CREATE TABLE departments  (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    name VARCHAR(30) NOT NULL
)
    PARTITION BY KEY(id)
    PARTITIONS 2;

INSERT INTO departments VALUES
    ('', 'Sales'), ('', 'Customer Service'),
    ('', 'Delivery'), ('', 'Accounting');
```

You can explicitly select partitions (or subpartitions, or both) from any or all of the tables in a join. (The `PARTITION` option used to select partitions from a given table immediately follows the name of the table, before all other options, including any table alias.) For example, the following query gets the name, employee ID, department, and city of all employees who work in the Sales or Delivery department (partition `p1` of the `departments` table) at the stores in either of the cities of Nambucca and Bellingen (partition `p0` of the `stores` table):

```
mysql> SELECT
    ->     e.id AS 'Employee ID', CONCAT(e.fname, ' ', e.lname) AS Name,
    ->     s.city AS City, d.name AS department
    -> FROM employees AS e
    ->     JOIN stores PARTITION (p1) AS s ON e.store_id=s.id
    ->     JOIN departments PARTITION (p0) AS d ON e.department_id=d.id
    -> ORDER BY e.lname;
+-------------+---------------+-----------+------------+
| Employee ID | Name          | City      | department |
+-------------+---------------+-----------+------------+
|          14 | Fred Goldberg | Bellingen | Delivery   |
|           5 | Mary Jones    | Nambucca  | Sales      |
|          17 | Mark Morgan   | Bellingen | Delivery   |
|           9 | Andy Smith    | Nambucca  | Delivery   |
|           8 | June Wilson   | Bellingen | Sales      |
+-------------+---------------+-----------+------------+
5 rows in set (0.00 sec)
```

For general information about joins in MySQL, see Section 13.2.10.2, “JOIN Syntax”.

When the `PARTITION` option is used with `DELETE` statements, only those partitions (and subpartitions, if any) listed with the option are checked for rows to be deleted. Any other partitions are ignored, as shown here:

```
mysql> SELECT * FROM employees WHERE fname LIKE 'j%';
+----+-------+--------+----------+---------------+
| id | fname | lname  | store_id | department_id |
+----+-------+--------+----------+---------------+
```

```
|  4 | Jim   | Smith  |        2 |             4 |
|  8 | June  | Wilson |        3 |             1 |
| 11 | Jill  | Stone  |        1 |             4 |
+----+-------+--------+----------+---------------+
3 rows in set (0.00 sec)

mysql> DELETE FROM employees PARTITION (p0, p1)
    ->     WHERE fname LIKE 'j%';
Query OK, 2 rows affected (0.09 sec)

mysql> SELECT * FROM employees WHERE fname LIKE 'j%';
+----+-------+-------+----------+---------------+
| id | fname | lname | store_id | department_id |
+----+-------+-------+----------+---------------+
| 11 | Jill  | Stone |        1 |             4 |
+----+-------+-------+----------+---------------+
1 row in set (0.00 sec)
```

Only the two rows in partitions `p0` and `p1` matching the `WHERE` condition were deleted. As you can see from the result when the `SELECT` is run a second time, there remains a row in the table matching the `WHERE` condition, but residing in a different partition (`p2`).

`UPDATE` statements using explicit partition selection behave in the same way; only rows in the partitions referenced by the `PARTITION` option are considered when determining the rows to be updated, as can be seen by executing the following statements:

```
mysql> UPDATE employees PARTITION (p0)
    ->     SET store_id = 2 WHERE fname = 'Jill';
Query OK, 0 rows affected (0.00 sec)
Rows matched: 0  Changed: 0  Warnings: 0

mysql> SELECT * FROM employees WHERE fname = 'Jill';
+----+-------+-------+----------+---------------+
| id | fname | lname | store_id | department_id |
+----+-------+-------+----------+---------------+
| 11 | Jill  | Stone |        1 |             4 |
+----+-------+-------+----------+---------------+
1 row in set (0.00 sec)

mysql> UPDATE employees PARTITION (p2)
    ->     SET store_id = 2 WHERE fname = 'Jill';
Query OK, 1 row affected (0.09 sec)
Rows matched: 1  Changed: 1  Warnings: 0

mysql> SELECT * FROM employees WHERE fname = 'Jill';
+----+-------+-------+----------+---------------+
| id | fname | lname | store_id | department_id |
+----+-------+-------+----------+---------------+
| 11 | Jill  | Stone |        2 |             4 |
+----+-------+-------+----------+---------------+
1 row in set (0.00 sec)
```

In the same way, when `PARTITION` is used with `DELETE`, only rows in the partition or partitions named in the partition list are checked for deletion.

For statements that insert rows, the behavior differs in that failure to find a suitable partition causes the statement to fail. This is true for both `INSERT` and `REPLACE` statements, as shown here:

```
mysql> INSERT INTO employees PARTITION (p2) VALUES (20, 'Jan', 'Jones', 1, 3);
ERROR 1729 (HY000): Found a row not matching the given partition set
mysql> INSERT INTO employees PARTITION (p3) VALUES (20, 'Jan', 'Jones', 1, 3);
Query OK, 1 row affected (0.07 sec)

mysql> REPLACE INTO employees PARTITION (p0) VALUES (20, 'Jan', 'Jones', 3, 2);
ERROR 1729 (HY000): Found a row not matching the given partition set

mysql> REPLACE INTO employees PARTITION (p3) VALUES (20, 'Jan', 'Jones', 3, 2);
Query OK, 2 rows affected (0.09 sec)
```

For statements that write multiple rows to a partitioned table that using the `InnoDB` storage engine: If any row in the list following `VALUES` cannot be written to one of the partitions specified in the *`partition_names`* list, the entire statement fails and no rows are written. This is shown for `INSERT` statements in the following example, reusing the `employees` table created previously:

```
mysql> ALTER TABLE employees
    ->     REORGANIZE PARTITION p3 INTO (
    ->         PARTITION p3 VALUES LESS THAN (20),
    ->         PARTITION p4 VALUES LESS THAN (25),
    ->         PARTITION p5 VALUES LESS THAN MAXVALUE
    ->     );
Query OK, 6 rows affected (2.09 sec)
Records: 6  Duplicates: 0  Warnings: 0

mysql> SHOW CREATE TABLE employees\G
*************************** 1. row ***************************
       Table: employees
Create Table: CREATE TABLE `employees` (
  `id` int(11) NOT NULL AUTO_INCREMENT,
  `fname` varchar(25) NOT NULL,
  `lname` varchar(25) NOT NULL,
  `store_id` int(11) NOT NULL,
  `department_id` int(11) NOT NULL,
  PRIMARY KEY (`id`)
) ENGINE=InnoDB AUTO_INCREMENT=27 DEFAULT CHARSET=utf8mb4
/*!50100 PARTITION BY RANGE (id)
(PARTITION p0 VALUES LESS THAN (5) ENGINE = InnoDB,
 PARTITION p1 VALUES LESS THAN (10) ENGINE = InnoDB,
 PARTITION p2 VALUES LESS THAN (15) ENGINE = InnoDB,
 PARTITION p3 VALUES LESS THAN (20) ENGINE = InnoDB,
 PARTITION p4 VALUES LESS THAN (25) ENGINE = InnoDB,
 PARTITION p5 VALUES LESS THAN MAXVALUE ENGINE = InnoDB) */
1 row in set (0.00 sec)

mysql> INSERT INTO employees PARTITION (p3, p4) VALUES
    ->     (24, 'Tim', 'Greene', 3, 1),  (26, 'Linda', 'Mills', 2, 1);
ERROR 1729 (HY000): Found a row not matching the given partition set

mysql> INSERT INTO employees PARTITION (p3, p4. p5) VALUES
    ->     (24, 'Tim', 'Greene', 3, 1),  (26, 'Linda', 'Mills', 2, 1);
Query OK, 2 rows affected (0.06 sec)
Records: 2  Duplicates: 0  Warnings: 0
```

The preceding is true for both `INSERT` statements and `REPLACE` statements that write multiple rows.

Partition selection is disabled for tables employing a storage engine that supplies automatic partitioning, such as `NDB` (used with MySQL Cluster; not currently supported in MySQL 8.0).

22.6 Restrictions and Limitations on Partitioning

This section discusses current restrictions and limitations on MySQL partitioning support.

Prohibited constructs. The following constructs are not permitted in partitioning expressions:

- Stored procedures, stored functions, UDFs, or plugins.
- Declared variables or user variables.

For a list of SQL functions which are permitted in partitioning expressions, see Section 22.6.3, “Partitioning Limitations Relating to Functions”.

Arithmetic and logical operators. Use of the arithmetic operators `+`, `-`, and `*` is permitted in partitioning expressions. However, the result must be an integer value or `NULL` (except in the case of `[LINEAR] KEY` partitioning, as discussed elsewhere in this chapter; see Section 22.2, “Partitioning Types”, for more information).

The `DIV` operator is also supported; the `/` operator is not permitted.

The bit operators `|`, `&`, `^`, `<<`, `>>`, and `~` are not permitted in partitioning expressions.

Server SQL mode. Tables employing user-defined partitioning do not preserve the SQL mode in effect at the time that they were created. As discussed elsewhere in this Manual (see Section 5.1.8, “Server SQL Modes”), the results of many MySQL functions and operators may change according to the server SQL mode. Therefore, a change in the SQL mode at any time after the creation of partitioned tables may lead to major changes in the behavior of such tables, and could easily lead to corruption or loss of data. For these reasons, *it is strongly recommended that you never change the server SQL mode after creating partitioned tables*.

Examples. The following examples illustrate some changes in behavior of partitioned tables due to a change in the server SQL mode:

1. **Error handling.** As discussed elsewhere, handling of “special” values such as zero and `NULL` can differ between different server SQL modes (see Section 5.1.8, “Server SQL Modes”). For example, `ERROR_FOR_DIVISION_BY_ZERO` can affect whether or not 0 can be inserted as a value into a table whose paritioning expression uses *`column DIV value`* or *`column MOD value`*.

2. **Table accessibility.** Sometimes a change in the server SQL mode can make partitioned tables unusable. The following `CREATE TABLE` statement can be executed successfully only if the `NO_UNSIGNED_SUBTRACTION` mode is in effect:

```
mysql> SELECT @@sql_mode;
+------------+
| @@sql_mode |
+------------+
|            |
+------------+
1 row in set (0.00 sec)

mysql> CREATE TABLE tu (c1 BIGINT UNSIGNED)
    ->   PARTITION BY RANGE(c1 - 10) (
    ->     PARTITION p0 VALUES LESS THAN (-5),
    ->     PARTITION p1 VALUES LESS THAN (0),
    ->     PARTITION p2 VALUES LESS THAN (5),
    ->     PARTITION p3 VALUES LESS THAN (10),
    ->     PARTITION p4 VALUES LESS THAN (MAXVALUE)
    -> );
ERROR 1563 (HY000): Partition constant is out of partition function domain

mysql> SET sql_mode='NO_UNSIGNED_SUBTRACTION';
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @@sql_mode;
+-------------------------+
| @@sql_mode              |
+-------------------------+
| NO_UNSIGNED_SUBTRACTION |
+-------------------------+
1 row in set (0.00 sec)

mysql> CREATE TABLE tu (c1 BIGINT UNSIGNED)
    ->   PARTITION BY RANGE(c1 - 10) (
    ->     PARTITION p0 VALUES LESS THAN (-5),
    ->     PARTITION p1 VALUES LESS THAN (0),
    ->     PARTITION p2 VALUES LESS THAN (5),
    ->     PARTITION p3 VALUES LESS THAN (10),
    ->     PARTITION p4 VALUES LESS THAN (MAXVALUE)
    -> );
Query OK, 0 rows affected (0.05 sec)
```

   If you remove the `NO_UNSIGNED_SUBTRACTION` server SQL mode after creating `tu`, you may no longer be able to access this table:

```
mysql> SET sql_mode='';
Query OK, 0 rows affected (0.00 sec)
```

```
mysql> SELECT * FROM tu;
ERROR 1563 (HY000): Partition constant is out of partition function domain
mysql> INSERT INTO tu VALUES (20);
ERROR 1563 (HY000): Partition constant is out of partition function domain
```

See also Section 5.1.8, “Server SQL Modes”.

Server SQL modes also impact replication of partitioned tables. Disparate SQL modes on master and slave can lead to partitioning expressions being evaluated differently; this can cause the distribution of data among partitions to be different in the master's and slave's copies of a given table, and may even cause inserts into partitioned tables that succeed on the master to fail on the slave. For best results, you should always use the same server SQL mode on the master and on the slave.

Performance considerations. Some effects of partitioning operations on performance are given in the following list:

- **File system operations.** Partitioning and repartitioning operations (such as `ALTER TABLE` with `PARTITION BY ...`, `REORGANIZE PARTITION`, or `REMOVE PARTITIONING`) depend on file system operations for their implementation. This means that the speed of these operations is affected by such factors as file system type and characteristics, disk speed, swap space, file handling efficiency of the operating system, and MySQL server options and variables that relate to file handling. In particular, you should make sure that `large_files_support` is enabled and that `open_files_limit` is set properly. Partitioning and repartitioning operations involving `InnoDB` tables may be made more efficient by enabling `innodb_file_per_table`.

  See also Maximum number of partitions.

- **Table locks.** Generally, the process executing a partitioning operation on a table takes a write lock on the table. Reads from such tables are relatively unaffected; pending `INSERT` and `UPDATE` operations are performed as soon as the partitioning operation has completed. For `InnoDB`-specific exceptions to this limitation, see Section 15.12.7, “Online DDL for Partitioned Tables”.

- **Indexes; partition pruning.** As with nonpartitioned tables, proper use of indexes can speed up queries on partitioned tables significantly. In addition, designing partitioned tables and queries on these tables to take advantage of *partition pruning* can improve performance dramatically. See Section 22.4, “Partition Pruning”, for more information.

  Index condition pushdown is supported for partitioned tables. See Section 8.2.1.5, “Index Condition Pushdown Optimization”.

- **Performance with LOAD DATA.** In MySQL 8.0, `LOAD DATA` uses buffering to improve performance. You should be aware that the buffer uses 130 KB memory per partition to achieve this.

Maximum number of partitions.
In MySQL 8.0, the maximum possible number of partitions for a given table is 8192. This number includes subpartitions.

If, when creating tables with a large number of partitions (but less than the maximum), you encounter an error message such as `Got error ... from storage engine: Out of resources when opening file`, you may be able to address the issue by increasing the value of the `open_files_limit` system variable. However, this is dependent on the operating system, and may not be possible or advisable on all platforms; see Section B.5.2.17, “File Not Found and Similar Errors”, for more information. In some cases, using large numbers (hundreds) of partitions may also not be advisable due to other concerns, so using more partitions does not automatically lead to better results.

See also File system operations.

Query cache not supported.
The query cache is not supported for partitioned tables, and is automatically disabled for queries involving partitioned tables. The query cache cannot be enabled for such queries.

Foreign keys not supported for partitioned InnoDB tables.
Partitioned tables using the `InnoDB` storage engine do not support foreign keys. More specifically, this means that the following two statements are true:

1. No definition of an `InnoDB` table employing user-defined partitioning may contain foreign key references; no `InnoDB` table whose definition contains foreign key references may be partitioned.

2. No `InnoDB` table definition may contain a foreign key reference to a user-partitioned table; no `InnoDB` table with user-defined partitioning may contain columns referenced by foreign keys.

The scope of the restrictions just listed includes all tables that use the `InnoDB` storage engine. `CREATE TABLE` and `ALTER TABLE` statements that would result in tables violating these restrictions are not allowed.

ALTER TABLE ... ORDER BY. An `ALTER TABLE ... ORDER BY` *`column`* statement run against a partitioned table causes ordering of rows only within each partition.

Effects on REPLACE statements by modification of primary keys. It can be desirable in some cases (see Section 22.6.1, “Partitioning Keys, Primary Keys, and Unique Keys”) to modify a table's primary key. Be aware that, if your application uses `REPLACE` statements and you do this, the results of these statements can be drastically altered. See Section 13.2.9, “REPLACE Syntax”, for more information and an example.

FULLTEXT indexes.
Partitioned tables do not support `FULLTEXT` indexes or searches.

Spatial columns. Columns with spatial data types such as `POINT` or `GEOMETRY` cannot be used in partitioned tables.

Temporary tables.
Temporary tables cannot be partitioned.

Log tables. It is not possible to partition the log tables; an `ALTER TABLE ... PARTITION BY ...` statement on such a table fails with an error.

Data type of partitioning key.
A partitioning key must be either an integer column or an expression that resolves to an integer. Expressions employing `ENUM` columns cannot be used. The column or expression value may also be `NULL`; see Section 22.2.7, “How MySQL Partitioning Handles NULL”.

There are two exceptions to this restriction:

1. When partitioning by [`LINEAR`] `KEY`, it is possible to use columns of any valid MySQL data type other than `TEXT` or `BLOB` as partitioning keys, because the internal key-hashing functions produce the correct data type from these types. For example, the following two `CREATE TABLE` statements are valid:

```
CREATE TABLE tkc (c1 CHAR)
PARTITION BY KEY(c1)
PARTITIONS 4;

CREATE TABLE tke
    ( c1 ENUM('red', 'orange', 'yellow', 'green', 'blue', 'indigo', 'violet') )
PARTITION BY LINEAR KEY(c1)
PARTITIONS 6;
```

2. When partitioning by `RANGE COLUMNS` or `LIST COLUMNS`, it is possible to use string, `DATE`, and `DATETIME` columns. For example, each of the following `CREATE TABLE` statements is valid:

```
CREATE TABLE rc (c1 INT, c2 DATE)
PARTITION BY RANGE COLUMNS(c2) (
```

```
    PARTITION p0 VALUES LESS THAN('1990-01-01'),
    PARTITION p1 VALUES LESS THAN('1995-01-01'),
    PARTITION p2 VALUES LESS THAN('2000-01-01'),
    PARTITION p3 VALUES LESS THAN('2005-01-01'),
    PARTITION p4 VALUES LESS THAN(MAXVALUE)
);

CREATE TABLE lc (c1 INT, c2 CHAR(1))
PARTITION BY LIST COLUMNS(c2) (
    PARTITION p0 VALUES IN('a', 'd', 'g', 'j', 'm', 'p', 's', 'v', 'y'),
    PARTITION p1 VALUES IN('b', 'e', 'h', 'k', 'n', 'q', 't', 'w', 'z'),
    PARTITION p2 VALUES IN('c', 'f', 'i', 'l', 'o', 'r', 'u', 'x', NULL)
);
```

Neither of the preceding exceptions applies to `BLOB` or `TEXT` column types.

Subqueries.
A partitioning key may not be a subquery, even if that subquery resolves to an integer value or `NULL`.

Issues with subpartitions.
Subpartitions must use `HASH` or `KEY` partitioning. Only `RANGE` and `LIST` partitions may be subpartitioned; `HASH` and `KEY` partitions cannot be subpartitioned.

`SUBPARTITION BY KEY` requires that the subpartitioning column or columns be specified explicitly, unlike the case with `PARTITION BY KEY`, where it can be omitted (in which case the table's primary key column is used by default). Consider the table created by this statement:

```
CREATE TABLE ts (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    name VARCHAR(30)
);
```

You can create a table having the same columns, partitioned by `KEY`, using a statement such as this one:

```
CREATE TABLE ts (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    name VARCHAR(30)
)
PARTITION BY KEY()
PARTITIONS 4;
```

The previous statement is treated as though it had been written like this, with the table's primary key column used as the partitioning column:

```
CREATE TABLE ts (
    id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    name VARCHAR(30)
)
PARTITION BY KEY(id)
PARTITIONS 4;
```

However, the following statement that attempts to create a subpartitioned table using the default column as the subpartitioning column fails, and the column must be specified for the statement to succeed, as shown here:

```
mysql> CREATE TABLE ts (
    ->     id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    ->     name VARCHAR(30)
    -> )
    -> PARTITION BY RANGE(id)
    -> SUBPARTITION BY KEY()
    -> SUBPARTITIONS 4
```

```
    -> (
    ->     PARTITION p0 VALUES LESS THAN (100),
    ->     PARTITION p1 VALUES LESS THAN (MAXVALUE)
    -> );
ERROR 1064 (42000): You have an error in your SQL syntax; check the manual that
corresponds to your MySQL server version for the right syntax to use near ')

mysql> CREATE TABLE ts (
    ->     id INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    ->     name VARCHAR(30)
    -> )
    -> PARTITION BY RANGE(id)
    -> SUBPARTITION BY KEY(id)
    -> SUBPARTITIONS 4
    -> (
    ->     PARTITION p0 VALUES LESS THAN (100),
    ->     PARTITION p1 VALUES LESS THAN (MAXVALUE)
    -> );
Query OK, 0 rows affected (0.07 sec)
```

This is a known issue (see Bug #51470).

DATA DIRECTORY and INDEX DIRECTORY options. Table-level `DATA DIRECTORY` and `INDEX DIRECTORY` options are ignored (see Bug #32091). You can employ these options for individual partitions or subpartitions of `InnoDB` tables.

Repairing and rebuilding partitioned tables. The statements `CHECK TABLE`, `OPTIMIZE TABLE`, `ANALYZE TABLE`, and `REPAIR TABLE` are supported for partitioned tables.

In addition, you can use `ALTER TABLE ... REBUILD PARTITION` to rebuild one or more partitions of a partitioned table; `ALTER TABLE ... REORGANIZE PARTITION` also causes partitions to be rebuilt. See Section 13.1.7, “ALTER TABLE Syntax”, for more information about these two statements.

`ANALYZE`, `CHECK`, `OPTIMIZE`, `REPAIR`, and `TRUNCATE` operations are supported with subpartitions. See Section 13.1.7.1, “ALTER TABLE Partition Operations”.

22.6.1 Partitioning Keys, Primary Keys, and Unique Keys

This section discusses the relationship of partitioning keys with primary keys and unique keys. The rule governing this relationship can be expressed as follows: All columns used in the partitioning expression for a partitioned table must be part of every unique key that the table may have.

In other words, *every unique key on the table must use every column in the table's partitioning expression*. (This also includes the table's primary key, since it is by definition a unique key. This particular case is discussed later in this section.) For example, each of the following table creation statements is invalid:

```
CREATE TABLE t1 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    UNIQUE KEY (col1, col2)
)
PARTITION BY HASH(col3)
PARTITIONS 4;

CREATE TABLE t2 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    UNIQUE KEY (col1),
    UNIQUE KEY (col3)
)
PARTITION BY HASH(col1 + col3)
```

```
PARTITIONS 4;
```

In each case, the proposed table would have at least one unique key that does not include all columns used in the partitioning expression.

Each of the following statements is valid, and represents one way in which the corresponding invalid table creation statement could be made to work:

```
CREATE TABLE t1 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    UNIQUE KEY (col1, col2, col3)
)
PARTITION BY HASH(col3)
PARTITIONS 4;

CREATE TABLE t2 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    UNIQUE KEY (col1, col3)
)
PARTITION BY HASH(col1 + col3)
PARTITIONS 4;
```

This example shows the error produced in such cases:

```
mysql> CREATE TABLE t3 (
    ->     col1 INT NOT NULL,
    ->     col2 DATE NOT NULL,
    ->     col3 INT NOT NULL,
    ->     col4 INT NOT NULL,
    ->     UNIQUE KEY (col1, col2),
    ->     UNIQUE KEY (col3)
    -> )
    -> PARTITION BY HASH(col1 + col3)
    -> PARTITIONS 4;
ERROR 1491 (HY000): A PRIMARY KEY must include all columns in the table's partitioning function
```

The `CREATE TABLE` statement fails because both `col1` and `col3` are included in the proposed partitioning key, but neither of these columns is part of both of unique keys on the table. This shows one possible fix for the invalid table definition:

```
mysql> CREATE TABLE t3 (
    ->     col1 INT NOT NULL,
    ->     col2 DATE NOT NULL,
    ->     col3 INT NOT NULL,
    ->     col4 INT NOT NULL,
    ->     UNIQUE KEY (col1, col2, col3),
    ->     UNIQUE KEY (col3)
    -> )
    -> PARTITION BY HASH(col3)
    -> PARTITIONS 4;
Query OK, 0 rows affected (0.05 sec)
```

In this case, the proposed partitioning key `col3` is part of both unique keys, and the table creation statement succeeds.

The following table cannot be partitioned at all, because there is no way to include in a partitioning key any columns that belong to both unique keys:

```
CREATE TABLE t4 (
```

```
    col1 INT NOT NULL,
    col2 INT NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    UNIQUE KEY (col1, col3),
    UNIQUE KEY (col2, col4)
);
```

Since every primary key is by definition a unique key, this restriction also includes the table's primary key, if it has one. For example, the next two statements are invalid:

```
CREATE TABLE t5 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    PRIMARY KEY(col1, col2)
)
PARTITION BY HASH(col3)
PARTITIONS 4;

CREATE TABLE t6 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    PRIMARY KEY(col1, col3),
    UNIQUE KEY(col2)
)
PARTITION BY HASH( YEAR(col2) )
PARTITIONS 4;
```

In both cases, the primary key does not include all columns referenced in the partitioning expression. However, both of the next two statements are valid:

```
CREATE TABLE t7 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    PRIMARY KEY(col1, col2)
)
PARTITION BY HASH(col1 + YEAR(col2))
PARTITIONS 4;

CREATE TABLE t8 (
    col1 INT NOT NULL,
    col2 DATE NOT NULL,
    col3 INT NOT NULL,
    col4 INT NOT NULL,
    PRIMARY KEY(col1, col2, col4),
    UNIQUE KEY(col2, col1)
)
PARTITION BY HASH(col1 + YEAR(col2))
PARTITIONS 4;
```

If a table has no unique keys—this includes having no primary key—then this restriction does not apply, and you may use any column or columns in the partitioning expression as long as the column type is compatible with the partitioning type.

For the same reason, you cannot later add a unique key to a partitioned table unless the key includes all columns used by the table's partitioning expression. Consider the partitioned table created as shown here:

```
mysql> CREATE TABLE t_no_pk (c1 INT, c2 INT)
    ->     PARTITION BY RANGE(c1) (
    ->         PARTITION p0 VALUES LESS THAN (10),
```

```
    ->         PARTITION p1 VALUES LESS THAN (20),
    ->         PARTITION p2 VALUES LESS THAN (30),
    ->         PARTITION p3 VALUES LESS THAN (40)
    ->     );
Query OK, 0 rows affected (0.12 sec)
```

It is possible to add a primary key to `t_no_pk` using either of these `ALTER TABLE` statements:

```
#  possible PK
mysql> ALTER TABLE t_no_pk ADD PRIMARY KEY(c1);
Query OK, 0 rows affected (0.13 sec)
Records: 0  Duplicates: 0  Warnings: 0

# drop this PK
mysql> ALTER TABLE t_no_pk DROP PRIMARY KEY;
Query OK, 0 rows affected (0.10 sec)
Records: 0  Duplicates: 0  Warnings: 0

#  use another possible PK
mysql> ALTER TABLE t_no_pk ADD PRIMARY KEY(c1, c2);
Query OK, 0 rows affected (0.12 sec)
Records: 0  Duplicates: 0  Warnings: 0

# drop this PK
mysql> ALTER TABLE t_no_pk DROP PRIMARY KEY;
Query OK, 0 rows affected (0.09 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

However, the next statement fails, because `c1` is part of the partitioning key, but is not part of the proposed primary key:

```
#  fails with error 1503
mysql> ALTER TABLE t_no_pk ADD PRIMARY KEY(c2);
ERROR 1503 (HY000): A PRIMARY KEY must include all columns in the table's partitioning function
```

Since `t_no_pk` has only `c1` in its partitioning expression, attempting to adding a unique key on `c2` alone fails. However, you can add a unique key that uses both `c1` and `c2`.

These rules also apply to existing nonpartitioned tables that you wish to partition using `ALTER TABLE ... PARTITION BY`. Consider a table `np_pk` created as shown here:

```
mysql> CREATE TABLE np_pk (
    ->     id INT NOT NULL AUTO_INCREMENT,
    ->     name VARCHAR(50),
    ->     added DATE,
    ->     PRIMARY KEY (id)
    -> );
Query OK, 0 rows affected (0.08 sec)
```

The following `ALTER TABLE` statement fails with an error, because the `added` column is not part of any unique key in the table:

```
mysql> ALTER TABLE np_pk
    ->     PARTITION BY HASH( TO_DAYS(added) )
    ->     PARTITIONS 4;
ERROR 1503 (HY000): A PRIMARY KEY must include all columns in the table's partitioning function
```

However, this statement using the `id` column for the partitioning column is valid, as shown here:

```
mysql> ALTER TABLE np_pk
    ->     PARTITION BY HASH(id)
    ->     PARTITIONS 4;
Query OK, 0 rows affected (0.11 sec)
Records: 0  Duplicates: 0  Warnings: 0
```

In the case of `np_pk`, the only column that may be used as part of a partitioning expression is `id`; if you wish to partition this table using any other column or columns in the partitioning expression, you must first modify the table, either by adding the desired column or columns to the primary key, or by dropping the primary key altogether.

22.6.2 Partitioning Limitations Relating to Storage Engines

In MySQL 8.0, partitioning support is not actually provided by the MySQL Server, but rather by a table storage engine's own or native partitioning handler. In MySQL 8.0, only the `InnoDB` storage engine provides a native partitioning handler. This means that partitioned tables cannot be created using any other storage engine.

Note

MySQL Cluster's `NDB` storage engine also provides native partitioning support, but is not currently supported in MySQL 8.0.

`ALTER TABLE ... OPTIMIZE PARTITION` does not work correctly with partitioned tables that use `InnoDB`. Use `ALTER TABLE ... REBUILD PARTITION` and `ALTER TABLE ... ANALYZE PARTITION`, instead, for such tables. For more information, see Section 13.1.7.1, "ALTER TABLE Partition Operations".

Upgrading partitioned tables. When performing an upgrade, tables which are partitioned by `KEY` must be dumped and reloaded. Partitioned tables using storage engines other than `InnoDB` cannot be upgraded from MySQL 5.7 or earlier to MySQL 8.0 or later; you must either drop the partitioning from such tables with `ALTER TABLE ... REMOVE PARTITIONING` or convert them to `InnoDB` using `ALTER TABLE ... ENGINE=INNODB` prior to the upgrade.

For information about converting `MyISAM` tables to `InnoDB`, see Section 15.8.1.4, "Converting Tables from MyISAM to InnoDB".

22.6.3 Partitioning Limitations Relating to Functions

This section discusses limitations in MySQL Partitioning relating specifically to functions used in partitioning expressions.

Only the MySQL functions shown in the following table are allowed in partitioning expressions.

`ABS()`	`CEILING()` (see CEILING() and FLOOR())	`DAY()`
`DAYOFMONTH()`	`DAYOFWEEK()`	`DAYOFYEAR()`
`DATEDIFF()`	`EXTRACT()` (see EXTRACT() function with WEEK specifier)	`FLOOR()` (see CEILING() and FLOOR())
`HOUR()`	`MICROSECOND()`	`MINUTE()`
`MOD()`	`MONTH()`	`QUARTER()`
`SECOND()`	`TIME_TO_SEC()`	`TO_DAYS()`
`TO_SECONDS()`	`UNIX_TIMESTAMP()` (with `TIMESTAMP` columns)	`WEEKDAY()`
`YEAR()`		`YEARWEEK()`

In MySQL 8.0, partition pruning is supported for the `TO_DAYS()`, `TO_SECONDS()`, `YEAR()`, and `UNIX_TIMESTAMP()` functions. See Section 22.4, "Partition Pruning", for more information.

CEILING() and FLOOR(). Each of these functions returns an integer only if it is passed an argument of an exact numeric type, such as one of the `INT` types or `DECIMAL`. This means, for example, that the following `CREATE TABLE` statement fails with an error, as shown here:

```
mysql> CREATE TABLE t (c FLOAT) PARTITION BY LIST( FLOOR(c) )(
    ->     PARTITION p0 VALUES IN (1,3,5),
    ->     PARTITION p1 VALUES IN (2,4,6)
    -> );
ERROR 1490 (HY000): The PARTITION function returns the wrong type
```

EXTRACT() function with WEEK specifier. The value returned by the `EXTRACT()` function, when used as `EXTRACT(WEEK FROM col)`, depends on the value of the `default_week_format` system variable. For this reason, `EXTRACT()` is not permitted as a partitioning function when it specifies the unit as `WEEK`. (Bug #54483)

See Section 12.6.2, “Mathematical Functions”, for more information about the return types of these functions, as well as Section 11.2, “Numeric Types”.

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This chapter discusses stored programs and views, which are database objects defined in terms of SQL code that is stored on the server for later execution.

Stored programs include these objects:

- Stored routines, that is, stored procedures and functions. A stored procedure is invoked using the `CALL` statement. A procedure does not have a return value but can modify its parameters for later inspection by the caller. It can also generate result sets to be returned to the client program. A stored function is used much like a built-in function. you invoke it in an expression and it returns a value during expression evaluation.
- Triggers. A trigger is a named database object that is associated with a table and that is activated when a particular event occurs for the table, such as an insert or update.
- Events. An event is a task that the server runs according to schedule.

Views are stored queries that when referenced produce a result set. A view acts as a virtual table.

This chapter describes how to use stored programs and views. The following sections provide additional information about SQL syntax for statements related to these objects:

- For each object type, there are `CREATE`, `ALTER`, and `DROP` statements that control which objects exist and how they are defined. See Section 13.1, “Data Definition Statements”.
- The `CALL` statement is used to invoke stored procedures. See Section 13.2.1, “CALL Syntax”.
- Stored program definitions include a body that may use compound statements, loops, conditionals, and declared variables. See Section 13.6, “Compound-Statement Syntax”.

In MySQL, metadata changes to objects referred to by stored programs are detected and cause automatic reparsing of the affected statements when the program is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

23.1 Defining Stored Programs

Each stored program contains a body that consists of an SQL statement. This statement may be a compound statement made up of several statements separated by semicolon (`;`) characters. For example, the following stored procedure has a body made up of a `BEGIN ... END` block that contains a `SET` statement and a `REPEAT` loop that itself contains another `SET` statement:

```
CREATE PROCEDURE dorepeat(p1 INT)
BEGIN
  SET @x = 0;
  REPEAT SET @x = @x + 1; UNTIL @x > p1 END REPEAT;
END;
```

If you use the `mysql` client program to define a stored program containing semicolon characters, a problem arises. By default, `mysql` itself recognizes the semicolon as a statement delimiter, so you must redefine the delimiter temporarily to cause `mysql` to pass the entire stored program definition to the server.

To redefine the `mysql` delimiter, use the `delimiter` command. The following example shows how to do this for the `dorepeat()` procedure just shown. The delimiter is changed to `//` to enable the entire definition to be passed to the server as a single statement, and then restored to `;` before invoking the procedure. This enables the `;` delimiter used in the procedure body to be passed through to the server rather than being interpreted by `mysql` itself.

```
mysql> delimiter //

mysql> CREATE PROCEDURE dorepeat(p1 INT)
    -> BEGIN
    ->   SET @x = 0;
    ->   REPEAT SET @x = @x + 1; UNTIL @x > p1 END REPEAT;
    -> END
    -> //
Query OK, 0 rows affected (0.00 sec)

mysql> delimiter ;

mysql> CALL dorepeat(1000);
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT @x;
+------+
| @x   |
+------+
| 1001 |
+------+
1 row in set (0.00 sec)
```

You can redefine the delimiter to a string other than `//`, and the delimiter can consist of a single character or multiple characters. You should avoid the use of the backslash (`\`) character because that is the escape character for MySQL.

The following is an example of a function that takes a parameter, performs an operation using an SQL function, and returns the result. In this case, it is unnecessary to use `delimiter` because the function definition contains no internal `;` statement delimiters:

```
mysql> CREATE FUNCTION hello (s CHAR(20))
mysql> RETURNS CHAR(50) DETERMINISTIC
    -> RETURN CONCAT('Hello, ',s,'!');
```

```
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT hello('world');
+----------------+
| hello('world') |
+----------------+
| Hello, world!  |
+----------------+
1 row in set (0.00 sec)
```

23.2 Using Stored Routines (Procedures and Functions)

MySQL supports stored routines (procedures and functions). A stored routine is a set of SQL statements that can be stored in the server. Once this has been done, clients don't need to keep reissuing the individual statements but can refer to the stored routine instead.

Stored routines can be particularly useful in certain situations:

- When multiple client applications are written in different languages or work on different platforms, but need to perform the same database operations.
- When security is paramount. Banks, for example, use stored procedures and functions for all common operations. This provides a consistent and secure environment, and routines can ensure that each operation is properly logged. In such a setup, applications and users would have no access to the database tables directly, but can only execute specific stored routines.

Stored routines can provide improved performance because less information needs to be sent between the server and the client. The tradeoff is that this does increase the load on the database server because more of the work is done on the server side and less is done on the client (application) side. Consider this if many client machines (such as Web servers) are serviced by only one or a few database servers.

Stored routines also enable you to have libraries of functions in the database server. This is a feature shared by modern application languages that enable such design internally (for example, by using classes). Using these client application language features is beneficial for the programmer even outside the scope of database use.

MySQL follows the SQL:2003 syntax for stored routines, which is also used by IBM's DB2. All syntax described here is supported and any limitations and extensions are documented where appropriate.

Additional Resources

- You may find the Stored Procedures User Forum of use when working with stored procedures and functions.
- For answers to some commonly asked questions regarding stored routines in MySQL, see Section A.4, “MySQL 8.0 FAQ: Stored Procedures and Functions”.
- There are some restrictions on the use of stored routines. See Section C.1, “Restrictions on Stored Programs”.
- Binary logging for stored routines takes place as described in Section 23.7, “Binary Logging of Stored Programs”.

23.2.1 Stored Routine Syntax

A stored routine is either a procedure or a function. Stored routines are created with the `CREATE PROCEDURE` and `CREATE FUNCTION` statements (see Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”). A procedure is invoked using a `CALL` statement (see Section 13.2.1, “CALL Syntax”), and can only pass back values using output variables. A function can be called

from inside a statement just like any other function (that is, by invoking the function's name), and can return a scalar value. The body of a stored routine can use compound statements (see Section 13.6, “Compound-Statement Syntax”).

Stored routines can be dropped with the `DROP PROCEDURE` and `DROP FUNCTION` statements (see Section 13.1.23, “DROP PROCEDURE and DROP FUNCTION Syntax”), and altered with the `ALTER PROCEDURE` and `ALTER FUNCTION` statements (see Section 13.1.5, “ALTER PROCEDURE Syntax”).

A stored procedure or function is associated with a particular database. This has several implications:

- When the routine is invoked, an implicit `USE db_name` is performed (and undone when the routine terminates). `USE` statements within stored routines are not permitted.
- You can qualify routine names with the database name. This can be used to refer to a routine that is not in the current database. For example, to invoke a stored procedure `p` or function `f` that is associated with the `test` database, you can say `CALL test.p()` or `test.f()`.
- When a database is dropped, all stored routines associated with it are dropped as well.

Stored functions cannot be recursive.

Recursion in stored procedures is permitted but disabled by default. To enable recursion, set the `max_sp_recursion_depth` server system variable to a value greater than zero. Stored procedure recursion increases the demand on thread stack space. If you increase the value of `max_sp_recursion_depth`, it may be necessary to increase thread stack size by increasing the value of `thread_stack` at server startup. See Section 5.1.5, “Server System Variables”, for more information.

MySQL supports a very useful extension that enables the use of regular `SELECT` statements (that is, without using cursors or local variables) inside a stored procedure. The result set of such a query is simply sent directly to the client. Multiple `SELECT` statements generate multiple result sets, so the client must use a MySQL client library that supports multiple result sets. This means the client must use a client library from a version of MySQL at least as recent as 4.1. The client should also specify the `CLIENT_MULTI_RESULTS` option when it connects. For C programs, this can be done with the `mysql_real_connect()` C API function. See Section 27.7.7.54, “mysql_real_connect()”, and Section 27.7.20, “C API Support for Multiple Statement Execution”.

23.2.2 Stored Routines and MySQL Privileges

The MySQL grant system takes stored routines into account as follows:

- The `CREATE ROUTINE` privilege is needed to create stored routines.
- The `ALTER ROUTINE` privilege is needed to alter or drop stored routines. This privilege is granted automatically to the creator of a routine if necessary, and dropped from the creator when the routine is dropped.
- The `EXECUTE` privilege is required to execute stored routines. However, this privilege is granted automatically to the creator of a routine if necessary (and dropped from the creator when the routine is dropped). Also, the default `SQL SECURITY` characteristic for a routine is `DEFINER`, which enables users who have access to the database with which the routine is associated to execute the routine.
- If the `automatic_sp_privileges` system variable is 0, the `EXECUTE` and `ALTER ROUTINE` privileges are not automatically granted to and dropped from the routine creator.
- The creator of a routine is the account used to execute the `CREATE` statement for it. This might not be the same as the account named as the `DEFINER` in the routine definition.

23.2.3 Stored Routine Metadata

Metadata about stored routines can be obtained as follows:

- Query the `ROUTINES` table of the `INFORMATION_SCHEMA` database. See Section 24.18, “The INFORMATION_SCHEMA ROUTINES Table”.
- Use the `SHOW CREATE PROCEDURE` and `SHOW CREATE FUNCTION` statements to see routine definitions. See Section 13.7.5.9, “SHOW CREATE PROCEDURE Syntax”.
- Use the `SHOW PROCEDURE STATUS` and `SHOW FUNCTION STATUS` statements to see routine characteristics. See Section 13.7.5.28, “SHOW PROCEDURE STATUS Syntax”.

23.2.4 Stored Procedures, Functions, Triggers, and LAST_INSERT_ID()

Within the body of a stored routine (procedure or function) or a trigger, the value of `LAST_INSERT_ID()` changes the same way as for statements executed outside the body of these kinds of objects (see Section 12.14, “Information Functions”). The effect of a stored routine or trigger upon the value of `LAST_INSERT_ID()` that is seen by following statements depends on the kind of routine:

- If a stored procedure executes statements that change the value of `LAST_INSERT_ID()`, the changed value is seen by statements that follow the procedure call.
- For stored functions and triggers that change the value, the value is restored when the function or trigger ends, so following statements do not see a changed value.

23.3 Using Triggers

A trigger is a named database object that is associated with a table, and that activates when a particular event occurs for the table. Some uses for triggers are to perform checks of values to be inserted into a table or to perform calculations on values involved in an update.

A trigger is defined to activate when a statement inserts, updates, or deletes rows in the associated table. These row operations are trigger events. For example, rows can be inserted by `INSERT` or `LOAD DATA` statements, and an insert trigger activates for each inserted row. A trigger can be set to activate either before or after the trigger event. For example, you can have a trigger activate before each row that is inserted into a table or after each row that is updated.

Important

MySQL triggers activate only for changes made to tables by SQL statements. This includes changes to base tables that underlie updatable views. Triggers do not activate for changes to tables made by APIs that do not transmit SQL statements to the MySQL Server. This means that triggers are not activated by updates made using the `NDB` API.

Triggers are not activated by changes in `INFORMATION_SCHEMA` or `performance_schema` tables. Those tables are actually views and triggers are not permitted on views.

The following sections describe the syntax for creating and dropping triggers, show some examples of how to use them, and indicate how to obtain trigger metadata.

Additional Resources

- You may find the Triggers User Forum of use when working with triggers.
- For answers to commonly asked questions regarding triggers in MySQL, see Section A.5, “MySQL 8.0 FAQ: Triggers”.

- There are some restrictions on the use of triggers; see Section C.1, “Restrictions on Stored Programs”.
- Binary logging for triggers takes place as described in Section 23.7, “Binary Logging of Stored Programs”.

23.3.1 Trigger Syntax and Examples

To create a trigger or drop a trigger, use the `CREATE TRIGGER` or `DROP TRIGGER` statement, described in Section 13.1.17, “CREATE TRIGGER Syntax”, and Section 13.1.27, “DROP TRIGGER Syntax”.

Here is a simple example that associates a trigger with a table, to activate for `INSERT` operations. The trigger acts as an accumulator, summing the values inserted into one of the columns of the table.

```
mysql> CREATE TABLE account (acct_num INT, amount DECIMAL(10,2));
Query OK, 0 rows affected (0.03 sec)

mysql> CREATE TRIGGER ins_sum BEFORE INSERT ON account
       FOR EACH ROW SET @sum = @sum + NEW.amount;
Query OK, 0 rows affected (0.01 sec)
```

The `CREATE TRIGGER` statement creates a trigger named `ins_sum` that is associated with the `account` table. It also includes clauses that specify the trigger action time, the triggering event, and what to do when the trigger activates:

- The keyword `BEFORE` indicates the trigger action time. In this case, the trigger activates before each row inserted into the table. The other permitted keyword here is `AFTER`.
- The keyword `INSERT` indicates the trigger event; that is, the type of operation that activates the trigger. In the example, `INSERT` operations cause trigger activation. You can also create triggers for `DELETE` and `UPDATE` operations.
- The statement following `FOR EACH ROW` defines the trigger body; that is, the statement to execute each time the trigger activates, which occurs once for each row affected by the triggering event. In the example, the trigger body is a simple `SET` that accumulates into a user variable the values inserted into the `amount` column. The statement refers to the column as `NEW.amount` which means “the value of the `amount` column to be inserted into the new row.”

To use the trigger, set the accumulator variable to zero, execute an `INSERT` statement, and then see what value the variable has afterward:

```
mysql> SET @sum = 0;
mysql> INSERT INTO account VALUES(137,14.98),(141,1937.50),(97,-100.00);
mysql> SELECT @sum AS 'Total amount inserted';
+-----------------------+
| Total amount inserted |
+-----------------------+
|               1852.48 |
+-----------------------+
```

In this case, the value of `@sum` after the `INSERT` statement has executed is `14.98 + 1937.50 - 100`, or `1852.48`.

To destroy the trigger, use a `DROP TRIGGER` statement. You must specify the schema name if the trigger is not in the default schema:

```
mysql> DROP TRIGGER test.ins_sum;
```

If you drop a table, any triggers for the table are also dropped.

Trigger names exist in the schema namespace, meaning that all triggers must have unique names within a schema. Triggers in different schemas can have the same name.

It is possible to define multiple triggers for a given table that have the same trigger event and action time. For example, you can have two `BEFORE UPDATE` triggers for a table. By default, triggers that have the same trigger event and action time activate in the order they were created. To affect trigger order, specify a clause after `FOR EACH ROW` that indicates `FOLLOWS` or `PRECEDES` and the name of an existing trigger that also has the same trigger event and action time. With `FOLLOWS`, the new trigger activates after the existing trigger. With `PRECEDES`, the new trigger activates before the existing trigger.

For example, the following trigger definition defines another `BEFORE INSERT` trigger for the `account` table:

```
mysql> CREATE TRIGGER ins_transaction BEFORE INSERT ON account
       FOR EACH ROW PRECEDES ins_sum
       SET
       @deposits = @deposits + IF(NEW.amount>0,NEW.amount,0),
       @withdrawals = @withdrawals + IF(NEW.amount<0,-NEW.amount,0);
Query OK, 0 rows affected (0.01 sec)
```

This trigger, `ins_transaction`, is similar to `ins_sum` but accumulates deposits and withdrawals separately. It has a `PRECEDES` clause that causes it to activate before `ins_sum`; without that clause, it would activate after `ins_sum` because it is created after `ins_sum`.

Within the trigger body, the `OLD` and `NEW` keywords enable you to access columns in the rows affected by a trigger. `OLD` and `NEW` are MySQL extensions to triggers; they are not case sensitive.

In an `INSERT` trigger, only `NEW.col_name` can be used; there is no old row. In a `DELETE` trigger, only `OLD.col_name` can be used; there is no new row. In an `UPDATE` trigger, you can use `OLD.col_name` to refer to the columns of a row before it is updated and `NEW.col_name` to refer to the columns of the row after it is updated.

A column named with `OLD` is read only. You can refer to it (if you have the `SELECT` privilege), but not modify it. You can refer to a column named with `NEW` if you have the `SELECT` privilege for it. In a `BEFORE` trigger, you can also change its value with `SET NEW.col_name = value` if you have the `UPDATE` privilege for it. This means you can use a trigger to modify the values to be inserted into a new row or used to update a row. (Such a `SET` statement has no effect in an `AFTER` trigger because the row change will have already occurred.)

In a `BEFORE` trigger, the `NEW` value for an `AUTO_INCREMENT` column is 0, not the sequence number that is generated automatically when the new row actually is inserted.

By using the `BEGIN ... END` construct, you can define a trigger that executes multiple statements. Within the `BEGIN` block, you also can use other syntax that is permitted within stored routines such as conditionals and loops. However, just as for stored routines, if you use the `mysql` program to define a trigger that executes multiple statements, it is necessary to redefine the `mysql` statement delimiter so that you can use the `;` statement delimiter within the trigger definition. The following example illustrates these points. It defines an `UPDATE` trigger that checks the new value to be used for updating each row, and modifies the value to be within the range from 0 to 100. This must be a `BEFORE` trigger because the value must be checked before it is used to update the row:

```
mysql> delimiter //
mysql> CREATE TRIGGER upd_check BEFORE UPDATE ON account
       FOR EACH ROW
       BEGIN
           IF NEW.amount < 0 THEN
               SET NEW.amount = 0;
           ELSEIF NEW.amount > 100 THEN
               SET NEW.amount = 100;
           END IF;
       END;//
```

```
mysql> delimiter ;
```

It can be easier to define a stored procedure separately and then invoke it from the trigger using a simple `CALL` statement. This is also advantageous if you want to execute the same code from within several triggers.

There are limitations on what can appear in statements that a trigger executes when activated:

- The trigger cannot use the `CALL` statement to invoke stored procedures that return data to the client or that use dynamic SQL. (Stored procedures are permitted to return data to the trigger through `OUT` or `INOUT` parameters.)
- The trigger cannot use statements that explicitly or implicitly begin or end a transaction, such as `START TRANSACTION`, `COMMIT`, or `ROLLBACK`. (`ROLLBACK to SAVEPOINT` is permitted because it does not end a transaction.).

See also Section C.1, “Restrictions on Stored Programs”.

MySQL handles errors during trigger execution as follows:

- If a `BEFORE` trigger fails, the operation on the corresponding row is not performed.
- A `BEFORE` trigger is activated by the *attempt* to insert or modify the row, regardless of whether the attempt subsequently succeeds.
- An `AFTER` trigger is executed only if any `BEFORE` triggers and the row operation execute successfully.
- An error during either a `BEFORE` or `AFTER` trigger results in failure of the entire statement that caused trigger invocation.
- For transactional tables, failure of a statement should cause rollback of all changes performed by the statement. Failure of a trigger causes the statement to fail, so trigger failure also causes rollback. For nontransactional tables, such rollback cannot be done, so although the statement fails, any changes performed prior to the point of the error remain in effect.

Triggers can contain direct references to tables by name, such as the trigger named `testref` shown in this example:

```
CREATE TABLE test1(a1 INT);
CREATE TABLE test2(a2 INT);
CREATE TABLE test3(a3 INT NOT NULL AUTO_INCREMENT PRIMARY KEY);
CREATE TABLE test4(
  a4 INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
  b4 INT DEFAULT 0
);

delimiter |

CREATE TRIGGER testref BEFORE INSERT ON test1
  FOR EACH ROW
  BEGIN
    INSERT INTO test2 SET a2 = NEW.a1;
    DELETE FROM test3 WHERE a3 = NEW.a1;
    UPDATE test4 SET b4 = b4 + 1 WHERE a4 = NEW.a1;
  END;
|

delimiter ;

INSERT INTO test3 (a3) VALUES
  (NULL), (NULL), (NULL), (NULL), (NULL),
  (NULL), (NULL), (NULL), (NULL), (NULL);

INSERT INTO test4 (a4) VALUES
  (0), (0), (0), (0), (0), (0), (0), (0), (0), (0);
```

Suppose that you insert the following values into table `test1` as shown here:

```
mysql> INSERT INTO test1 VALUES
       (1), (3), (1), (7), (1), (8), (4), (4);
Query OK, 8 rows affected (0.01 sec)
Records: 8  Duplicates: 0  Warnings: 0
```

As a result, the four tables contain the following data:

```
mysql> SELECT * FROM test1;
+------+
| a1   |
+------+
|    1 |
|    3 |
|    1 |
|    7 |
|    1 |
|    8 |
|    4 |
|    4 |
+------+
8 rows in set (0.00 sec)

mysql> SELECT * FROM test2;
+------+
| a2   |
+------+
|    1 |
|    3 |
|    1 |
|    7 |
|    1 |
|    8 |
|    4 |
|    4 |
+------+
8 rows in set (0.00 sec)

mysql> SELECT * FROM test3;
+----+
| a3 |
+----+
|  2 |
|  5 |
|  6 |
|  9 |
| 10 |
+----+
5 rows in set (0.00 sec)

mysql> SELECT * FROM test4;
+----+------+
| a4 | b4   |
+----+------+
|  1 |    3 |
|  2 |    0 |
|  3 |    1 |
|  4 |    2 |
|  5 |    0 |
|  6 |    0 |
|  7 |    1 |
|  8 |    1 |
|  9 |    0 |
| 10 |    0 |
+----+------+
10 rows in set (0.00 sec)
```

23.3.2 Trigger Metadata

Metadata about triggers can be obtained as follows:

- Query the `TRIGGERS` table of the `INFORMATION_SCHEMA` database. See Section 24.28, “The INFORMATION_SCHEMA TRIGGERS Table”.
- Use the `SHOW CREATE TRIGGER` statement. See Section 13.7.5.11, “SHOW CREATE TRIGGER Syntax”.
- Use the `SHOW TRIGGERS` statement. See Section 13.7.5.38, “SHOW TRIGGERS Syntax”.

23.4 Using the Event Scheduler

The *MySQL Event Scheduler* manages the scheduling and execution of events, that is, tasks that run according to a schedule. The following discussion covers the Event Scheduler and is divided into the following sections:

- Section 23.4.1, “Event Scheduler Overview”, provides an introduction to and conceptual overview of MySQL Events.
- Section 23.4.3, “Event Syntax”, discusses the SQL statements for creating, altering, and dropping MySQL Events.
- Section 23.4.4, “Event Metadata”, shows how to obtain information about events and how this information is stored by the MySQL Server.
- Section 23.4.6, “The Event Scheduler and MySQL Privileges”, discusses the privileges required to work with events and the ramifications that events have with regard to privileges when executing.

Stored routines require the `events` data dictionary table in the `mysql` system database. This table is created during the MySQL 8.0 installation procedure. If you are upgrading to MySQL 8.0 from an earlier version, be sure to run `mysql_upgrade` to to make sure that your system database is up to date. See Section 4.4.5, “`mysql_upgrade` — Check and Upgrade MySQL Tables”.

Additional Resources

- You may find the MySQL Event Scheduler User Forum of use when working with scheduled events.
- There are some restrictions on the use of events; see Section C.1, “Restrictions on Stored Programs”.
- Binary logging for events takes place as described in Section 23.7, “Binary Logging of Stored Programs”.

23.4.1 Event Scheduler Overview

MySQL Events are tasks that run according to a schedule. Therefore, we sometimes refer to them as *scheduled* events. When you create an event, you are creating a named database object containing one or more SQL statements to be executed at one or more regular intervals, beginning and ending at a specific date and time. Conceptually, this is similar to the idea of the Unix `crontab` (also known as a “cron job”) or the Windows Task Scheduler.

Scheduled tasks of this type are also sometimes known as “temporal triggers”, implying that these are objects that are triggered by the passage of time. While this is essentially correct, we prefer to use the term *events* to avoid confusion with triggers of the type discussed in Section 23.3, “Using Triggers”. Events should more specifically not be confused with “temporary triggers”. Whereas a trigger is a database object whose statements are executed in response to a specific type of event that occurs on a given table, a (scheduled) event is an object whose statements are executed in response to the passage of a specified time interval.

While there is no provision in the SQL Standard for event scheduling, there are precedents in other database systems, and you may notice some similarities between these implementations and that found in the MySQL Server.

MySQL Events have the following major features and properties:

- In MySQL, an event is uniquely identified by its name and the schema to which it is assigned.
- An event performs a specific action according to a schedule. This action consists of an SQL statement, which can be a compound statement in a `BEGIN ... END` block if desired (see Section 13.6, “Compound-Statement Syntax”). An event's timing can be either *one-time* or *recurrent*. A one-time event executes one time only. A recurrent event repeats its action at a regular interval, and the schedule for a recurring event can be assigned a specific start day and time, end day and time, both, or neither. (By default, a recurring event's schedule begins as soon as it is created, and continues indefinitely, until it is disabled or dropped.)

  If a repeating event does not terminate within its scheduling interval, the result may be multiple instances of the event executing simultaneously. If this is undesirable, you should institute a mechanism to prevent simultaneous instances. For example, you could use the `GET_LOCK()` function, or row or table locking.
- Users can create, modify, and drop scheduled events using SQL statements intended for these purposes. Syntactically invalid event creation and modification statements fail with an appropriate error message. *A user may include statements in an event's action which require privileges that the user does not actually have*. The event creation or modification statement succeeds but the event's action fails. See Section 23.4.6, “The Event Scheduler and MySQL Privileges” for details.
- Many of the properties of an event can be set or modified using SQL statements. These properties include the event's name, timing, persistence (that is, whether it is preserved following the expiration of its schedule), status (enabled or disabled), action to be performed, and the schema to which it is assigned. See Section 13.1.2, “ALTER EVENT Syntax”.

  The default definer of an event is the user who created the event, unless the event has been altered, in which case the definer is the user who issued the last `ALTER EVENT` statement affecting that event. An event can be modified by any user having the `EVENT` privilege on the database for which the event is defined. See Section 23.4.6, “The Event Scheduler and MySQL Privileges”.
- An event's action statement may include most SQL statements permitted within stored routines. For restrictions, see Section C.1, “Restrictions on Stored Programs”.

23.4.2 Event Scheduler Configuration

Events are executed by a special *event scheduler thread*; when we refer to the Event Scheduler, we actually refer to this thread. When running, the event scheduler thread and its current state can be seen by users having the `PROCESS` privilege in the output of `SHOW PROCESSLIST`, as shown in the discussion that follows.

The global `event_scheduler` system variable determines whether the Event Scheduler is enabled and running on the server. It has one of these 3 values, which affect event scheduling as described here:

- `OFF`: The Event Scheduler is stopped. The event scheduler thread does not run, is not shown in the output of `SHOW PROCESSLIST`, and no scheduled events are executed. `OFF` is the default value for `event_scheduler`.

  When the Event Scheduler is stopped (`event_scheduler` is `OFF`), it can be started by setting the value of `event_scheduler` to `ON`. (See next item.)
- `ON`: The Event Scheduler is started; the event scheduler thread runs and executes all scheduled events.

When the Event Scheduler is `ON`, the event scheduler thread is listed in the output of `SHOW PROCESSLIST` as a daemon process, and its state is represented as shown here:

```
mysql> SHOW PROCESSLIST\G
*************************** 1. row ***************************
     Id: 1
   User: root
   Host: localhost
     db: NULL
Command: Query
   Time: 0
  State: NULL
   Info: show processlist
*************************** 2. row ***************************
     Id: 2
   User: event_scheduler
   Host: localhost
     db: NULL
Command: Daemon
   Time: 3
  State: Waiting for next activation
   Info: NULL
2 rows in set (0.00 sec)
```

Event scheduling can be stopped by setting the value of `event_scheduler` to `OFF`.

- `DISABLED`: This value renders the Event Scheduler nonoperational. When the Event Scheduler is `DISABLED`, the event scheduler thread does not run (and so does not appear in the output of `SHOW PROCESSLIST`). In addition, the Event Scheduler state cannot be changed at runtime.

If the Event Scheduler status has not been set to `DISABLED`, `event_scheduler` can be toggled between `ON` and `OFF` (using `SET`). It is also possible to use `0` for `OFF`, and `1` for `ON` when setting this variable. Thus, any of the following 4 statements can be used in the `mysql` client to turn on the Event Scheduler:

```
SET GLOBAL event_scheduler = ON;
SET @@global.event_scheduler = ON;
SET GLOBAL event_scheduler = 1;
SET @@global.event_scheduler = 1;
```

Similarly, any of these 4 statements can be used to turn off the Event Scheduler:

```
SET GLOBAL event_scheduler = OFF;
SET @@global.event_scheduler = OFF;
SET GLOBAL event_scheduler = 0;
SET @@global.event_scheduler = 0;
```

Although `ON` and `OFF` have numeric equivalents, the value displayed for `event_scheduler` by `SELECT` or `SHOW VARIABLES` is always one of `OFF`, `ON`, or `DISABLED`. *`DISABLED` has no numeric equivalent*. For this reason, `ON` and `OFF` are usually preferred over `1` and `0` when setting this variable.

Note that attempting to set `event_scheduler` without specifying it as a global variable causes an error:

```
mysql< SET @@event_scheduler = OFF;
ERROR 1229 (HY000): Variable 'event_scheduler' is a GLOBAL
variable and should be set with SET GLOBAL
```

Important

It is possible to set the Event Scheduler to `DISABLED` only at server startup. If `event_scheduler` is `ON` or `OFF`, you cannot set it to `DISABLED` at runtime.

> Also, if the Event Scheduler is set to `DISABLED` at startup, you cannot change the value of `event_scheduler` at runtime.

To disable the event scheduler, use one of the following two methods:

- As a command-line option when starting the server:

```
--event-scheduler=DISABLED
```

- In the server configuration file (`my.cnf`, or `my.ini` on Windows systems), include the line where it will be read by the server (for example, in a `[mysqld]` section):

```
event_scheduler=DISABLED
```

To enable the Event Scheduler, restart the server without the `--event-scheduler=DISABLED` command-line option, or after removing or commenting out the line containing `event-scheduler=DISABLED` in the server configuration file, as appropriate. Alternatively, you can use `ON` (or `1`) or `OFF` (or `0`) in place of the `DISABLED` value when starting the server.

Note

You can issue event-manipulation statements when `event_scheduler` is set to `DISABLED`. No warnings or errors are generated in such cases (provided that the statements are themselves valid). However, scheduled events cannot execute until this variable is set to `ON` (or `1`). Once this has been done, the event scheduler thread executes all events whose scheduling conditions are satisfied.

Starting the MySQL server with the `--skip-grant-tables` option causes `event_scheduler` to be set to `DISABLED`, overriding any other value set either on the command line or in the `my.cnf` or `my.ini` file (Bug #26807).

For SQL statements used to create, alter, and drop events, see Section 23.4.3, “Event Syntax”.

MySQL provides an `EVENTS` table in the `INFORMATION_SCHEMA` database. This table can be queried to obtain information about scheduled events which have been defined on the server. See Section 23.4.4, “Event Metadata”, and Section 24.8, “The INFORMATION_SCHEMA EVENTS Table”, for more information.

For information regarding event scheduling and the MySQL privilege system, see Section 23.4.6, “The Event Scheduler and MySQL Privileges”.

23.4.3 Event Syntax

MySQL provides several SQL statements for working with scheduled events:

- New events are defined using the `CREATE EVENT` statement. See Section 13.1.10, “CREATE EVENT Syntax”.
- The definition of an existing event can be changed by means of the `ALTER EVENT` statement. See Section 13.1.2, “ALTER EVENT Syntax”.
- When a scheduled event is no longer wanted or needed, it can be deleted from the server by its definer using the `DROP EVENT` statement. See Section 13.1.20, “DROP EVENT Syntax”. Whether an event persists past the end of its schedule also depends on its `ON COMPLETION` clause, if it has one. See Section 13.1.10, “CREATE EVENT Syntax”.

  An event can be dropped by any user having the `EVENT` privilege for the database on which the event is defined. See Section 23.4.6, “The Event Scheduler and MySQL Privileges”.

23.4.4 Event Metadata

Metadata about events can be obtained as follows:

- Query the `EVENTS` table of the `INFORMATION_SCHEMA` database. See Section 24.8, "The INFORMATION_SCHEMA EVENTS Table".
- Use the `SHOW CREATE EVENT` statement. See Section 13.7.5.7, "SHOW CREATE EVENT Syntax".
- Use the `SHOW EVENTS` statement. See Section 13.7.5.18, "SHOW EVENTS Syntax".

Event Scheduler Time Representation

Each session in MySQL has a session time zone (STZ). This is the session `time_zone` value that is initialized from the server's global `time_zone` value when the session begins but may be changed during the session.

The session time zone that is current when a `CREATE EVENT` or `ALTER EVENT` statement executes is used to interpret times specified in the event definition. This becomes the event time zone (ETZ); that is, the time zone that is used for event scheduling and is in effect within the event as it executes.

For representation of event information in the `mysql.event` table, the `execute_at`, `starts`, and `ends` times are converted to UTC and stored along with the event time zone. This enables event execution to proceed as defined regardless of any subsequent changes to the server time zone or daylight saving time effects. The `last_executed` time is also stored in UTC.

If you select information from `mysql.event`, the times just mentioned are retrieved as UTC values. These times can also be obtained by selecting from the `INFORMATION_SCHEMA.EVENTS` table or from `SHOW EVENTS`, but they are reported as ETZ values. Other times available from these sources indicate when an event was created or last altered; these are displayed as STZ values. The following table summarizes representation of event times.

Value	`mysql.event`	`INFORMATION_SCHEMA.EVENTS`	`SHOW EVENTS`
Execute at	UTC	ETZ	ETZ
Starts	UTC	ETZ	ETZ
Ends	UTC	ETZ	ETZ
Last executed	UTC	ETZ	n/a
Created	STZ	STZ	n/a
Last altered	STZ	STZ	n/a

23.4.5 Event Scheduler Status

The Event Scheduler writes information about event execution that terminates with an error or warning to the MySQL Server's error log. See Section 23.4.6, "The Event Scheduler and MySQL Privileges" for an example.

To obtain information about the state of the Event Scheduler for debugging and troubleshooting purposes, run `mysqladmin debug` (see Section 4.5.2, "`mysqladmin` — Client for Administering a MySQL Server"); after running this command, the server's error log contains output relating to the Event Scheduler, similar to what is shown here:

```
Events status:
LLA = Last Locked At  LUA = Last Unlocked At
WOC = Waiting On Condition  DL = Data Locked

Event scheduler status:
State      : INITIALIZED
Thread id  : 0
LLA        : init_scheduler:313
```

```
LUA        : init_scheduler:318
WOC        : NO
Workers    : 0
Executed   : 0
Data locked: NO

Event queue status:
Element count   : 1
Data locked     : NO
Attempting lock : NO
LLA             : init_queue:148
LUA             : init_queue:168
WOC             : NO
Next activation : 0000-00-00 00:00:00
```

In statements that occur as part of events executed by the Event Scheduler, diagnostics messages (not only errors, but also warnings) are written to the error log, and, on Windows, to the application event log. For frequently executed events, it is possible for this to result in many logged messages. For example, for `SELECT ... INTO var_list` statements, if the query returns no rows, a warning with error code 1329 occurs (`No data`), and the variable values remain unchanged. If the query returns multiple rows, error 1172 occurs (`Result consisted of more than one row`). For either condition, you can avoid having the warnings be logged by declaring a condition handler; see Section 13.6.7.2, “DECLARE ... HANDLER Syntax”. For statements that may retrieve multiple rows, another strategy is to use `LIMIT 1` to limit the result set to a single row.

23.4.6 The Event Scheduler and MySQL Privileges

To enable or disable the execution of scheduled events, it is necessary to set the value of the global `event_scheduler` system variable. This requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege.

The `EVENT` privilege governs the creation, modification, and deletion of events. This privilege can be bestowed using `GRANT`. For example, this `GRANT` statement confers the `EVENT` privilege for the schema named `myschema` on the user `jon@ghidora`:

```
GRANT EVENT ON myschema.* TO jon@ghidora;
```

(We assume that this user account already exists, and that we wish for it to remain unchanged otherwise.)

To grant this same user the `EVENT` privilege on all schemas, use the following statement:

```
GRANT EVENT ON *.* TO jon@ghidora;
```

The `EVENT` privilege has global or schema-level scope. Therefore, trying to grant it on a single table results in an error as shown:

```
mysql> GRANT EVENT ON myschema.mytable TO jon@ghidora;
ERROR 1144 (42000): Illegal GRANT/REVOKE command; please
consult the manual to see which privileges can be used
```

It is important to understand that an event is executed with the privileges of its definer, and that it cannot perform any actions for which its definer does not have the requisite privileges. For example, suppose that `jon@ghidora` has the `EVENT` privilege for `myschema`. Suppose also that this user has the `SELECT` privilege for `myschema`, but no other privileges for this schema. It is possible for `jon@ghidora` to create a new event such as this one:

```
CREATE EVENT e_store_ts
    ON SCHEDULE
      EVERY 10 SECOND
    DO
```

```
      INSERT INTO myschema.mytable VALUES (UNIX_TIMESTAMP());
```

The user waits for a minute or so, and then performs a `SELECT * FROM mytable;` query, expecting to see several new rows in the table. Instead, the table is empty. Since the user does not have the `INSERT` privilege for the table in question, the event has no effect.

If you inspect the MySQL error log (*`hostname`*`.err`), you can see that the event is executing, but the action it is attempting to perform fails:

```
2013-09-24T12:41:31.261992Z 25 [ERROR] Event Scheduler:
[jon@ghidora][cookbook.e_store_ts] INSERT command denied to user
'jon'@'ghidora' for table 'mytable'
2013-09-24T12:41:31.262022Z 25 [Note] Event Scheduler:
[jon@ghidora].[myschema.e_store_ts] event execution failed.
2013-09-24T12:41:41.271796Z 26 [ERROR] Event Scheduler:
[jon@ghidora][cookbook.e_store_ts] INSERT command denied to user
'jon'@'ghidora' for table 'mytable'
2013-09-24T12:41:41.272761Z 26 [Note] Event Scheduler:
[jon@ghidora].[myschema.e_store_ts] event execution failed.
```

Since this user very likely does not have access to the error log, it is possible to verify whether the event's action statement is valid by executing it directly:

```
mysql> INSERT INTO myschema.mytable VALUES (UNIX_TIMESTAMP());
ERROR 1142 (42000): INSERT command denied to user
'jon'@'ghidora' for table 'mytable'
```

Inspection of the `INFORMATION_SCHEMA.EVENTS` table shows that `e_store_ts` exists and is enabled, but its `LAST_EXECUTED` column is `NULL`:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.EVENTS
     >     WHERE EVENT_NAME='e_store_ts'
     >     AND EVENT_SCHEMA='myschema'\G
*************************** 1. row ***************************
   EVENT_CATALOG: NULL
    EVENT_SCHEMA: myschema
      EVENT_NAME: e_store_ts
         DEFINER: jon@ghidora
      EVENT_BODY: SQL
EVENT_DEFINITION: INSERT INTO myschema.mytable VALUES (UNIX_TIMESTAMP())
      EVENT_TYPE: RECURRING
      EXECUTE_AT: NULL
  INTERVAL_VALUE: 5
  INTERVAL_FIELD: SECOND
        SQL_MODE: NULL
          STARTS: 0000-00-00 00:00:00
            ENDS: 0000-00-00 00:00:00
          STATUS: ENABLED
   ON_COMPLETION: NOT PRESERVE
         CREATED: 2006-02-09 22:36:06
    LAST_ALTERED: 2006-02-09 22:36:06
   LAST_EXECUTED: NULL
   EVENT_COMMENT:
1 row in set (0.00 sec)
```

To rescind the `EVENT` privilege, use the `REVOKE` statement. In this example, the `EVENT` privilege on the schema `myschema` is removed from the `jon@ghidora` user account:

```
REVOKE EVENT ON myschema.* FROM jon@ghidora;
```

Important

Revoking the `EVENT` privilege from a user does not delete or disable any events that may have been created by that user.

> An event is not migrated or dropped as a result of renaming or dropping the user who created it.

Suppose that the user `jon@ghidora` has been granted the `EVENT` and `INSERT` privileges on the `myschema` schema. This user then creates the following event:

```
CREATE EVENT e_insert
    ON SCHEDULE
      EVERY 7 SECOND
    DO
      INSERT INTO myschema.mytable;
```

After this event has been created, `root` revokes the `EVENT` privilege for `jon@ghidora`. However, `e_insert` continues to execute, inserting a new row into `mytable` each seven seconds. The same would be true if `root` had issued either of these statements:

- `DROP USER jon@ghidora;`
- `RENAME USER jon@ghidora TO someotherguy@ghidora;`

You can verify that this is true by examining the `mysql.event` table (discussed later in this section) or the `INFORMATION_SCHEMA.EVENTS` table (see Section 24.8, “The INFORMATION_SCHEMA EVENTS Table”) before and after issuing a `DROP USER` or `RENAME USER` statement.

Event definitions are stored in the `mysql.event` table. To drop an event created by another user account, the MySQL `root` user (or another user with the necessary privileges) can delete rows from this table. For example, to remove the event `e_insert` shown previously, `root` can use the following statement:

```
DELETE FROM mysql.event
    WHERE db = 'myschema'
      AND definer = 'jon@ghidora'
      AND name = 'e_insert';
```

It is very important to match the event name, database schema name, and user account when deleting rows from the `mysql.event` table. This is because the same user can create different events of the same name in different schemas.

Users' `EVENT` privileges are stored in the `Event_priv` columns of the `mysql.user` and `mysql.db` tables. In both cases, this column holds one of the values '`Y`' or '`N`'. '`N`' is the default. `mysql.user.Event_priv` is set to '`Y`' for a given user only if that user has the global `EVENT` privilege (that is, if the privilege was bestowed using `GRANT EVENT ON *.*`). For a schema-level `EVENT` privilege, `GRANT` creates a row in `mysql.db` and sets that row's `Db` column to the name of the schema, the `User` column to the name of the user, and the `Event_priv` column to '`Y`'. There should never be any need to manipulate these tables directly, since the `GRANT EVENT` and `REVOKE EVENT` statements perform the required operations on them.

Five status variables provide counts of event-related operations (but *not* of statements executed by events; see Section C.1, “Restrictions on Stored Programs”). These are:

- `Com_create_event`: The number of `CREATE EVENT` statements executed since the last server restart.
- `Com_alter_event`: The number of `ALTER EVENT` statements executed since the last server restart.
- `Com_drop_event`: The number of `DROP EVENT` statements executed since the last server restart.
- `Com_show_create_event`: The number of `SHOW CREATE EVENT` statements executed since the last server restart.

- `Com_show_events`: The number of `SHOW EVENTS` statements executed since the last server restart.

You can view current values for all of these at one time by running the statement `SHOW STATUS LIKE '%event%';`.

23.5 Using Views

MySQL supports views, including updatable views. Views are stored queries that when invoked produce a result set. A view acts as a virtual table.

The following discussion describes the syntax for creating and dropping views, and shows some examples of how to use them.

Additional Resources

- You may find the Views User Forum of use when working with views.
- For answers to some commonly asked questions regarding views in MySQL, see Section A.6, “MySQL 8.0 FAQ: Views”.
- There are some restrictions on the use of views; see Section C.5, “Restrictions on Views”.

23.5.1 View Syntax

The `CREATE VIEW` statement creates a new view (see Section 13.1.18, “CREATE VIEW Syntax”). To alter the definition of a view or drop a view, use `ALTER VIEW` (see Section 13.1.8, “ALTER VIEW Syntax”), or `DROP VIEW` (see Section 13.1.28, “DROP VIEW Syntax”).

A view can be created from many kinds of `SELECT` statements. It can refer to base tables or other views. It can use joins, `UNION`, and subqueries. The `SELECT` need not even refer to any tables. The following example defines a view that selects two columns from another table, as well as an expression calculated from those columns:

```
mysql> CREATE TABLE t (qty INT, price INT);
mysql> INSERT INTO t VALUES(3, 50), (5, 60);
mysql> CREATE VIEW v AS SELECT qty, price, qty*price AS value FROM t;
mysql> SELECT * FROM v;
+------+-------+-------+
| qty  | price | value |
+------+-------+-------+
|    3 |    50 |   150 |
|    5 |    60 |   300 |
+------+-------+-------+
mysql> SELECT * FROM v WHERE qty = 5;
+------+-------+-------+
| qty  | price | value |
+------+-------+-------+
|    5 |    60 |   300 |
+------+-------+-------+
```

23.5.2 View Processing Algorithms

The optional `ALGORITHM` clause for `CREATE VIEW` or `ALTER VIEW` is a MySQL extension to standard SQL. It affects how MySQL processes the view. `ALGORITHM` takes three values: `MERGE`, `TEMPTABLE`, or `UNDEFINED`.

- For `MERGE`, the text of a statement that refers to the view and the view definition are merged such that parts of the view definition replace corresponding parts of the statement.
- For `TEMPTABLE`, the results from the view are retrieved into a temporary table, which then is used to execute the statement.

- For `UNDEFINED`, MySQL chooses which algorithm to use. It prefers `MERGE` over `TEMPTABLE` if possible, because `MERGE` is usually more efficient and because a view cannot be updatable if a temporary table is used.
- If no `ALGORITHM` clause is present, the default algorithm is determined by the value of the `derived_merge` flag of the `optimizer_switch` system variable. For additional discussion, see Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

A reason to specify `TEMPTABLE` explicitly is that locks can be released on underlying tables after the temporary table has been created and before it is used to finish processing the statement. This might result in quicker lock release than the `MERGE` algorithm so that other clients that use the view are not blocked as long.

A view algorithm can be `UNDEFINED` for three reasons:

- No `ALGORITHM` clause is present in the `CREATE VIEW` statement.
- The `CREATE VIEW` statement has an explicit `ALGORITHM = UNDEFINED` clause.
- `ALGORITHM = MERGE` is specified for a view that can be processed only with a temporary table. In this case, MySQL generates a warning and sets the algorithm to `UNDEFINED`.

As mentioned earlier, `MERGE` is handled by merging corresponding parts of a view definition into the statement that refers to the view. The following examples briefly illustrate how the `MERGE` algorithm works. The examples assume that there is a view `v_merge` that has this definition:

```
CREATE ALGORITHM = MERGE VIEW v_merge (vc1, vc2) AS
SELECT c1, c2 FROM t WHERE c3 > 100;
```

Example 1: Suppose that we issue this statement:

```
SELECT * FROM v_merge;
```

MySQL handles the statement as follows:

- `v_merge` becomes `t`
- `*` becomes `vc1, vc2`, which corresponds to `c1, c2`
- The view `WHERE` clause is added

The resulting statement to be executed becomes:

```
SELECT c1, c2 FROM t WHERE c3 > 100;
```

Example 2: Suppose that we issue this statement:

```
SELECT * FROM v_merge WHERE vc1 < 100;
```

This statement is handled similarly to the previous one, except that `vc1 < 100` becomes `c1 < 100` and the view `WHERE` clause is added to the statement `WHERE` clause using an `AND` connective (and parentheses are added to make sure the parts of the clause are executed with correct precedence). The resulting statement to be executed becomes:

```
SELECT c1, c2 FROM t WHERE (c3 > 100) AND (c1 < 100);
```

Effectively, the statement to be executed has a `WHERE` clause of this form:

```
```

```
WHERE (select WHERE) AND (view WHERE)
```

If the MERGE algorithm cannot be used, a temporary table must be used instead. Constructs that prevent merging are the same as those that prevent merging in derived tables. Examples are SELECT DISTINCT or LIMIT in the subquery. For details, see Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

23.5.3 Updatable and Insertable Views

Some views are updatable and references to them can be used to specify tables to be updated in data change statements. That is, you can use them in statements such as UPDATE, DELETE, or INSERT to update the contents of the underlying table. Derived tables can also be specified in multiple-table UPDATE and DELETE statements, but can only be used for reading data to specify rows to be updated or deleted. Generally, the view references must be updatable, meaning that they may be merged and not materialized. Composite views have more complex rules.

For a view to be updatable, there must be a one-to-one relationship between the rows in the view and the rows in the underlying table. There are also certain other constructs that make a view nonupdatable. To be more specific, a view is not updatable if it contains any of the following:

- Aggregate functions (SUM(), MIN(), MAX(), COUNT(), and so forth)
- DISTINCT
- GROUP BY
- HAVING
- UNION or UNION ALL
- Subquery in the select list

  Nondependent subqueries in the select list fail for INSERT, but are okay for UPDATE, DELETE. For dependent subqueries in the select list, no data change statements are permitted.
- Certain joins (see additional join discussion later in this section)
- Reference to nonupdatable view in the FROM clause
- Subquery in the WHERE clause that refers to a table in the FROM clause
- Refers only to literal values (in this case, there is no underlying table to update)
- ALGORITHM = TEMPTABLE (use of a temporary table always makes a view nonupdatable)
- Multiple references to any column of a base table (fails for INSERT, okay for UPDATE, DELETE)

A generated column in a view is considered updatable because it is possible to assign to it. However, if such a column is updated explicitly, the only permitted value is DEFAULT. For information about generated columns, see Section 13.1.15.8, “CREATE TABLE and Generated Columns”.

It is sometimes possible for a multiple-table view to be updatable, assuming that it can be processed with the MERGE algorithm. For this to work, the view must use an inner join (not an outer join or a UNION). Also, only a single table in the view definition can be updated, so the SET clause must name only columns from one of the tables in the view. Views that use UNION ALL are not permitted even though they might be theoretically updatable.

With respect to insertability (being updatable with INSERT statements), an updatable view is insertable if it also satisfies these additional requirements for the view columns:

- There must be no duplicate view column names.

- The view must contain all columns in the base table that do not have a default value.
- The view columns must be simple column references. They must not be expressions, such as these:

```
3.14159
col1 + 3
UPPER(col2)
col3 / col4
(subquery)
```

MySQL sets a flag, called the view updatability flag, at `CREATE VIEW` time. The flag is set to `YES` (true) if `UPDATE` and `DELETE` (and similar operations) are legal for the view. Otherwise, the flag is set to `NO` (false). The `IS_UPDATABLE` column in the `INFORMATION_SCHEMA.VIEWS` table displays the status of this flag. It means that the server always knows whether a view is updatable.

If a view is not updatable, statements such `UPDATE`, `DELETE`, and `INSERT` are illegal and are rejected. (Note that even if a view is updatable, it might not be possible to insert into it, as described elsewhere in this section.)

The updatability of views may be affected by the value of the `updatable_views_with_limit` system variable. See Section 5.1.5, “Server System Variables”.

For the following discussion, suppose that these tables and views exist:

```
CREATE TABLE t1 (x INTEGER);
CREATE TABLE t2 (c INTEGER);
CREATE VIEW vmat AS SELECT SUM(x) AS s FROM t1;
CREATE VIEW vup AS SELECT * FROM t2;
CREATE VIEW vjoin AS SELECT * FROM vmat JOIN vup ON vmat.s=vup.c;
```

`INSERT`, `UPDATE`, and `DELETE` statements are permitted as follows:

- `INSERT`: The insert table of an `INSERT` statement may be a view reference that is merged. If the view is a join view, all components of the view must be updatable (not materialized). For a multiple-table updatable view, `INSERT` can work if it inserts into a single table.

  This statement is invalid because one component of the join view is nonupdatable:

  ```
  INSERT INTO vjoin (c) VALUES (1);
  ```

  This statement is valid; the view contains no materialized components:

  ```
  INSERT INTO vup (c) VALUES (1);
  ```

- `UPDATE`: The table or tables to be updated in an `UPDATE` statement may be view references that are merged. If a view is a join view, at least one component of the view must be updatable (this differs from `INSERT`).

  In a multiple-table `UPDATE` statement, the updated table references of the statement must be base tables or updatable view references. Nonupdated table references may be materialized views or derived tables.

  This statement is valid; column `c` is from the updatable part of the join view:

  ```
  UPDATE vjoin SET c=c+1;
  ```

  This statement is invalid; column `x` is from the nonupdatable part:

  ```
  UPDATE vjoin SET x=x+1;
  ```

This statement is valid; the updated table reference of the multiple-table `UPDATE` is an updatable view (`vup`):

```
UPDATE vup JOIN (SELECT SUM(x) AS s FROM t1) AS dt ON ...
SET c=c+1;
```

This statement is invalid; it tries to update a materialized derived table:

```
UPDATE vup JOIN (SELECT SUM(x) AS s FROM t1) AS dt ON ...
SET s=s+1;
```

- `DELETE`: The table or tables to be deleted from in a `DELETE` statement must be merged views. Join views are not allowed (this differs from `INSERT` and `UPDATE`).

  This statement is invalid because the view is a join view:

  ```
  DELETE vjoin WHERE ...;
  ```

  This statement is valid because the view is a merged (updatable) view:

  ```
  DELETE vup WHERE ...;
  ```

  This statement is valid because it deletes from a merged (updatable) view:

  ```
  DELETE vup FROM vup JOIN (SELECT SUM(x) AS s FROM t1) AS dt ON ...;
  ```

Additional discussion and examples follow.

Earlier discussion in this section pointed out that a view is not insertable if not all columns are simple column references (for example, if it contains columns that are expressions or composite expressions). Although such a view is not insertable, it can be updatable if you update only columns that are not expressions. Consider this view:

```
CREATE VIEW v AS SELECT col1, 1 AS col2 FROM t;
```

This view is not insertable because `col2` is an expression. But it is updatable if the update does not try to update `col2`. This update is permissible:

```
UPDATE v SET col1 = 0;
```

This update is not permissible because it attempts to update an expression column:

```
UPDATE v SET col2 = 0;
```

If a table contains an `AUTO_INCREMENT` column, inserting into an insertable view on the table that does not include the `AUTO_INCREMENT` column does not change the value of `LAST_INSERT_ID()`, because the side effects of inserting default values into columns not part of the view should not be visible.

23.5.4 The View WITH CHECK OPTION Clause

The `WITH CHECK OPTION` clause can be given for an updatable view to prevent inserts to rows for which the `WHERE` clause in the *`select_statement`* is not true. It also prevents updates to rows for which the `WHERE` clause is true but the update would cause it to be not true (in other words, it prevents visible rows from being updated to nonvisible rows).

In a WITH CHECK OPTION clause for an updatable view, the LOCAL and CASCADED keywords determine the scope of check testing when the view is defined in terms of another view. When neither keyword is given, the default is CASCADED.

WITH CHECK OPTION testing is standard-compliant:

- With LOCAL, the view WHERE clause is checked, then checking recurses to underlying views and applies the same rules.
- With CASCADED, the view WHERE clause is checked, then checking recurses to underlying views, adds WITH CASCADED CHECK OPTION to them (for purposes of the check; their definitions remain unchanged), and applies the same rules.
- With no check option, the view WHERE clause is not checked, then checking recurses to underlying views, and applies the same rules.

Consider the definitions for the following table and set of views:

```
CREATE TABLE t1 (a INT);
CREATE VIEW v1 AS SELECT * FROM t1 WHERE a < 2
WITH CHECK OPTION;
CREATE VIEW v2 AS SELECT * FROM v1 WHERE a > 0
WITH LOCAL CHECK OPTION;
CREATE VIEW v3 AS SELECT * FROM v1 WHERE a > 0
WITH CASCADED CHECK OPTION;
```

Here the v2 and v3 views are defined in terms of another view, v1.

Inserts for v2 are checked against its LOCAL check option, then the check recurses to v1 and the rules are applied again. The rules for v1 cause a check failure. The check for v3 also fails:

```
mysql> INSERT INTO v2 VALUES (2);
ERROR 1369 (HY000): CHECK OPTION failed 'test.v2'
mysql> INSERT INTO v3 VALUES (2);
ERROR 1369 (HY000): CHECK OPTION failed 'test.v3'
```

23.5.5 View Metadata

Metadata about views can be obtained as follows:

- Query the VIEWS table of the INFORMATION_SCHEMA database. See Section 24.30, “The INFORMATION_SCHEMA VIEWS Table”.
- Use the SHOW CREATE VIEW statement. See Section 13.7.5.13, “SHOW CREATE VIEW Syntax”.

23.6 Access Control for Stored Programs and Views

Stored programs and views are defined prior to use and, when referenced, execute within a security context that determines their privileges. These privileges are controlled by their DEFINER attribute, and, if there is one, their SQL SECURITY characteristic.

All stored programs (procedures, functions, triggers, and events) and views can have a DEFINER attribute that names a MySQL account. If the DEFINER attribute is omitted from a stored program or view definition, the default account is the user who creates the object.

In addition, stored routines (procedures and functions) and views can have an SQL SECURITY characteristic with a value of DEFINER or INVOKER to specify whether the object executes in definer or invoker context. If the SQL SECURITY characteristic is omitted, the default is definer context.

Triggers and events have no SQL SECURITY characteristic and always execute in definer context. The server invokes these objects automatically as necessary, so there is no invoking user.

Definer and invoker security contexts differ as follows:

- A stored program or view that executes in definer security context executes with the privileges of the account named by its `DEFINER` attribute. These privileges may be entirely different from those of the invoking user. The invoker must have appropriate privileges to reference the object (for example, `EXECUTE` to call a stored procedure or `SELECT` to select from a view), but when the object executes, the invoker's privileges are ignored and only the `DEFINER` account privileges matter. If this account has few privileges, the object is correspondingly limited in the operations it can perform. If the `DEFINER` account is highly privileged (such as a `root` account), the object can perform powerful operations *no matter who invokes it.*
- A stored routine or view that executes in invoker security context can perform only operations for which the invoker has privileges. The `DEFINER` attribute can be specified but has no effect for objects that execute in invoker context.

Consider the following stored procedure:

```
CREATE DEFINER = 'admin'@'localhost' PROCEDURE p1()
SQL SECURITY DEFINER
BEGIN
  UPDATE t1 SET counter = counter + 1;
END;
```

Any user who has the `EXECUTE` privilege for `p1` can invoke it with a `CALL` statement. However, when `p1` executes, it does so in `DEFINER` security context and thus executes with the privileges of `'admin'@'localhost'`, the account named in the `DEFINER` attribute. This account must have the `EXECUTE` privilege for `p1` as well as the `UPDATE` privilege for the table `t1`. Otherwise, the procedure fails.

Now consider this stored procedure, which is identical to `p1` except that its `SQL SECURITY` characteristic is `INVOKER`:

```
CREATE DEFINER = 'admin'@'localhost' PROCEDURE p2()
SQL SECURITY INVOKER
BEGIN
  UPDATE t1 SET counter = counter + 1;
END;
```

`p2`, unlike `p1`, executes in `INVOKER` security context. The `DEFINER` attribute is irrelevant and `p2` executes with the privileges of the invoking user. `p2` fails if the invoker lacks the `EXECUTE` privilege for `p2` or the `UPDATE` privilege for the table `t1`.

MySQL uses the following rules to control which accounts a user can specify in an object `DEFINER` attribute:

- You can specify a `DEFINER` value other than your own account only if you have the `SET_USER_ID` or `SUPER` privilege.
- If you do not have the `SET_USER_ID` or `SUPER` privilege, the only legal user value is your own account, either specified literally or by using `CURRENT_USER`. You cannot set the definer to some other account.

To minimize the risk potential for stored program and view creation and use, follow these guidelines:

- For a stored routine or view, use `SQL SECURITY INVOKER` in the object definition when possible so that it can be used only by users with permissions appropriate for the operations performed by the object.
- If you create definer-context stored programs or views while using an account that has the `SET_USER_ID` or `SUPER` privilege, specify an explicit `DEFINER` attribute that names an account

possessing only the privileges required for the operations performed by the object. Specify a highly privileged `DEFINER` account only when absolutely necessary.

- Administrators can prevent users from specifying highly privileged `DEFINER` accounts by not granting them the `SET_USER_ID` or `SUPER` privilege.
- Definer-context objects should be written keeping in mind that they may be able to access data for which the invoking user has no privileges. In some cases, you can prevent reference to these objects by not granting unauthorized users particular privileges:
  - A stored procedure or function cannot be referenced by a user who does not have the `EXECUTE` privilege for it.
  - A view cannot be referenced by a user who does not have the appropriate privilege for it (`SELECT` to select from it, `INSERT` to insert into it, and so forth).

  However, no such control exists for triggers because users do not reference them directly. A trigger always executes in `DEFINER` context and is activated by access to the table with which it is associated, even ordinary table accesses by users with no special privileges. If the `DEFINER` account is highly privileged, the trigger can perform sensitive or dangerous operations. This remains true if the `SET_USER_ID` (or `SUPER`) and `TRIGGER` privileges needed to create the trigger are revoked from the account of the user who created it. Administrators should be especially careful about granting users that combination of privileges.

23.7 Binary Logging of Stored Programs

The binary log contains information about SQL statements that modify database contents. This information is stored in the form of “events” that describe the modifications. The binary log has two important purposes:

- For replication, the binary log is used on master replication servers as a record of the statements to be sent to slave servers. The master server sends the events contained in its binary log to its slaves, which execute those events to make the same data changes that were made on the master. See Section 18.2, “Replication Implementation”.
- Certain data recovery operations require use of the binary log. After a backup file has been restored, the events in the binary log that were recorded after the backup was made are re-executed. These events bring databases up to date from the point of the backup. See Section 7.3.2, “Using Backups for Recovery”.

However, if logging occurs at the statement level, there are certain binary logging issues with respect to stored programs (stored procedures and functions, triggers, and events):

- In some cases, a statement might affect different sets of rows on master and slave.
- Replicated statements executed on a slave are processed by the slave SQL thread, which has full privileges. It is possible for a procedure to follow different execution paths on master and slave servers, so a user can write a routine containing a dangerous statement that will execute only on the slave where it is processed by a thread that has full privileges.
- If a stored program that modifies data is nondeterministic, it is not repeatable. This can result in different data on master and slave, or cause restored data to differ from the original data.

This section describes how MySQL handles binary logging for stored programs. It states the current conditions that the implementation places on the use of stored programs, and what you can do to avoid logging problems. It also provides additional information about the reasons for these conditions.

In general, the issues described here result when binary logging occurs at the SQL statement level. If you use row-based binary logging, the log contains changes made to individual rows as a result of executing SQL statements. When routines or triggers execute, row changes are logged, not the

statements that make the changes. For stored procedures, this means that the `CALL` statement is not logged. For stored functions, row changes made within the function are logged, not the function invocation. For triggers, row changes made by the trigger are logged. On the slave side, only the row changes are seen, not the stored program invocation. For general information about row-based logging, see Section 18.2.1, “Replication Formats”.

Unless noted otherwise, the remarks here assume that you have enabled binary logging by starting the server with the `--log-bin` option. (See Section 5.4.4, “The Binary Log”.) If the binary log is not enabled, replication is not possible, nor is the binary log available for data recovery.

The conditions on the use of stored functions in MySQL can be summarized as follows. These conditions do not apply to stored procedures or Event Scheduler events and they do not apply unless binary logging is enabled.

- To create or alter a stored function, you must have the `SET_USER_ID` or `SUPER` privilege, in addition to the `CREATE ROUTINE` or `ALTER ROUTINE` privilege that is normally required. (Depending on the `DEFINER` value in the function definition, `SET_USER_ID` or `SUPER` might be required regardless of whether binary logging is enabled. See Section 13.1.13, “CREATE PROCEDURE and CREATE FUNCTION Syntax”.)
- When you create a stored function, you must declare either that it is deterministic or that it does not modify data. Otherwise, it may be unsafe for data recovery or replication.

  By default, for a `CREATE FUNCTION` statement to be accepted, at least one of `DETERMINISTIC`, `NO SQL`, or `READS SQL DATA` must be specified explicitly. Otherwise an error occurs:

  ```
  ERROR 1418 (HY000): This function has none of DETERMINISTIC, NO SQL,
  or READS SQL DATA in its declaration and binary logging is enabled
  (you *might* want to use the less safe log_bin_trust_function_creators
  variable)
  ```

  This function is deterministic (and does not modify data), so it is safe:

  ```
  CREATE FUNCTION f1(i INT)
  RETURNS INT
  DETERMINISTIC
  READS SQL DATA
  BEGIN
    RETURN i;
  END;
  ```

  This function uses `UUID()`, which is not deterministic, so the function also is not deterministic and is not safe:

  ```
  CREATE FUNCTION f2()
  RETURNS CHAR(36) CHARACTER SET utf8
  BEGIN
    RETURN UUID();
  END;
  ```

  This function modifies data, so it may not be safe:

  ```
  CREATE FUNCTION f3(p_id INT)
  RETURNS INT
  BEGIN
    UPDATE t SET modtime = NOW() WHERE id = p_id;
    RETURN ROW_COUNT();
  END;
  ```

  Assessment of the nature of a function is based on the “honesty” of the creator: MySQL does not check that a function declared `DETERMINISTIC` is free of statements that produce nondeterministic results.

- Although it is possible to create a deterministic stored function without specifying `DETERMINISTIC`, you cannot execute this function using statement-based binary logging. To execute such a function, you must use row-based or mixed binary logging. Alternatively, if you explicitly specify `DETERMINISTIC` in the function definition, you can use any kind of logging, including statement-based logging.

- To relax the preceding conditions on function creation (that you must have the `SUPER` privilege and that a function must be declared deterministic or to not modify data), set the global `log_bin_trust_function_creators` system variable to 1. By default, this variable has a value of 0, but you can change it like this:

  ```
  mysql> SET GLOBAL log_bin_trust_function_creators = 1;
  ```

  You can also set this variable by using the `--log-bin-trust-function-creators=1` option when starting the server.

  If binary logging is not enabled, `log_bin_trust_function_creators` does not apply. `SUPER` is not required for function creation unless, as described previously, the `DEFINER` value in the function definition requires it.

- For information about built-in functions that may be unsafe for replication (and thus cause stored functions that use them to be unsafe as well), see Section 18.4.1, “Replication Features and Issues”.

Triggers are similar to stored functions, so the preceding remarks regarding functions also apply to triggers with the following exception: `CREATE TRIGGER` does not have an optional `DETERMINISTIC` characteristic, so triggers are assumed to be always deterministic. However, this assumption might be invalid in some cases. For example, the `UUID()` function is nondeterministic (and does not replicate). Be careful about using such functions in triggers.

Triggers can update tables, so error messages similar to those for stored functions occur with `CREATE TRIGGER` if you do not have the required privileges. On the slave side, the slave uses the trigger `DEFINER` attribute to determine which user is considered to be the creator of the trigger.

The rest of this section provides additional detail about the logging implementation and its implications. You need not read it unless you are interested in the background on the rationale for the current logging-related conditions on stored routine use. This discussion applies only for statement-based logging, and not for row-based logging, with the exception of the first item: `CREATE` and `DROP` statements are logged as statements regardless of the logging mode.

- The server writes `CREATE EVENT`, `CREATE PROCEDURE`, `CREATE FUNCTION`, `ALTER EVENT`, `ALTER PROCEDURE`, `ALTER FUNCTION`, `DROP EVENT`, `DROP PROCEDURE`, and `DROP FUNCTION` statements to the binary log.

- A stored function invocation is logged as a `SELECT` statement if the function changes data and occurs within a statement that would not otherwise be logged. This prevents nonreplication of data changes that result from use of stored functions in nonlogged statements. For example, `SELECT` statements are not written to the binary log, but a `SELECT` might invoke a stored function that makes changes. To handle this, a `SELECT` *`func_name`*`()` statement is written to the binary log when the given function makes a change. Suppose that the following statements are executed on the master:

  ```
  CREATE FUNCTION f1(a INT) RETURNS INT
  BEGIN
    IF (a < 3) THEN
      INSERT INTO t2 VALUES (a);
    END IF;
    RETURN 0;
  END;

  CREATE TABLE t1 (a INT);
  INSERT INTO t1 VALUES (1),(2),(3);
  ```

```
SELECT f1(a) FROM t1;
```

When the `SELECT` statement executes, the function `f1()` is invoked three times. Two of those invocations insert a row, and MySQL logs a `SELECT` statement for each of them. That is, MySQL writes the following statements to the binary log:

```
SELECT f1(1);
SELECT f1(2);
```

The server also logs a `SELECT` statement for a stored function invocation when the function invokes a stored procedure that causes an error. In this case, the server writes the `SELECT` statement to the log along with the expected error code. On the slave, if the same error occurs, that is the expected result and replication continues. Otherwise, replication stops.

- Logging stored function invocations rather than the statements executed by a function has a security implication for replication, which arises from two factors:
  - It is possible for a function to follow different execution paths on master and slave servers.
  - Statements executed on a slave are processed by the slave SQL thread which has full privileges.

  The implication is that although a user must have the `CREATE ROUTINE` privilege to create a function, the user can write a function containing a dangerous statement that will execute only on the slave where it is processed by a thread that has full privileges. For example, if the master and slave servers have server ID values of 1 and 2, respectively, a user on the master server could create and invoke an unsafe function `unsafe_func()` as follows:

```
mysql> delimiter //
mysql> CREATE FUNCTION unsafe_func () RETURNS INT
    -> BEGIN
    ->   IF @@server_id=2 THEN dangerous_statement; END IF;
    ->   RETURN 1;
    -> END;
    -> //
mysql> delimiter ;
mysql> INSERT INTO t VALUES(unsafe_func());
```

  The `CREATE FUNCTION` and `INSERT` statements are written to the binary log, so the slave will execute them. Because the slave SQL thread has full privileges, it will execute the dangerous statement. Thus, the function invocation has different effects on the master and slave and is not replication-safe.

  To guard against this danger for servers that have binary logging enabled, stored function creators must have the `SUPER` privilege, in addition to the usual `CREATE ROUTINE` privilege that is required. Similarly, to use `ALTER FUNCTION`, you must have the `SUPER` privilege in addition to the `ALTER ROUTINE` privilege. Without the `SUPER` privilege, an error will occur:

```
ERROR 1419 (HY000): You do not have the SUPER privilege and
binary logging is enabled (you *might* want to use the less safe
log_bin_trust_function_creators variable)
```

  If you do not want to require function creators to have the `SUPER` privilege (for example, if all users with the `CREATE ROUTINE` privilege on your system are experienced application developers), set the global `log_bin_trust_function_creators` system variable to 1. You can also set this variable by using the `--log-bin-trust-function-creators=1` option when starting the server. If binary logging is not enabled, `log_bin_trust_function_creators` does not apply. `SUPER` is not required for function creation unless, as described previously, the `DEFINER` value in the function definition requires it.

- If a function that performs updates is nondeterministic, it is not repeatable. This can have two undesirable effects:

- It will make a slave different from the master.
- Restored data will be different from the original data.

To deal with these problems, MySQL enforces the following requirement: On a master server, creation and alteration of a function is refused unless you declare the function to be deterministic or to not modify data. Two sets of function characteristics apply here:

- The `DETERMINISTIC` and `NOT DETERMINISTIC` characteristics indicate whether a function always produces the same result for given inputs. The default is `NOT DETERMINISTIC` if neither characteristic is given. To declare that a function is deterministic, you must specify `DETERMINISTIC` explicitly.
- The `CONTAINS SQL`, `NO SQL`, `READS SQL DATA`, and `MODIFIES SQL DATA` characteristics provide information about whether the function reads or writes data. Either `NO SQL` or `READS SQL DATA` indicates that a function does not change data, but you must specify one of these explicitly because the default is `CONTAINS SQL` if no characteristic is given.

By default, for a `CREATE FUNCTION` statement to be accepted, at least one of `DETERMINISTIC`, `NO SQL`, or `READS SQL DATA` must be specified explicitly. Otherwise an error occurs:

```
ERROR 1418 (HY000): This function has none of DETERMINISTIC, NO SQL,
or READS SQL DATA in its declaration and binary logging is enabled
(you *might* want to use the less safe log_bin_trust_function_creators
variable)
```

If you set `log_bin_trust_function_creators` to 1, the requirement that functions be deterministic or not modify data is dropped.

- Stored procedure calls are logged at the statement level rather than at the `CALL` level. That is, the server does not log the `CALL` statement, it logs those statements within the procedure that actually execute. As a result, the same changes that occur on the master will be observed on slave servers. This prevents problems that could result from a procedure having different execution paths on different machines.

  In general, statements executed within a stored procedure are written to the binary log using the same rules that would apply were the statements to be executed in standalone fashion. Some special care is taken when logging procedure statements because statement execution within procedures is not quite the same as in nonprocedure context:

  - A statement to be logged might contain references to local procedure variables. These variables do not exist outside of stored procedure context, so a statement that refers to such a variable cannot be logged literally. Instead, each reference to a local variable is replaced by this construct for logging purposes:

    ```
    NAME_CONST(var_name, var_value)
    ```

    `var_name` is the local variable name, and *`var_value`* is a constant indicating the value that the variable has at the time the statement is logged. `NAME_CONST()` has a value of *`var_value`*, and a “name” of *`var_name`*. Thus, if you invoke this function directly, you get a result like this:

    ```
    mysql> SELECT NAME_CONST('myname', 14);
    +--------+
    | myname |
    +--------+
    |     14 |
    +--------+
    ```

`NAME_CONST()` enables a logged standalone statement to be executed on a slave with the same effect as the original statement that was executed on the master within a stored procedure.

The use of `NAME_CONST()` can result in a problem for `CREATE TABLE ... SELECT` statements when the source column expressions refer to local variables. Converting these references to `NAME_CONST()` expressions can result in column names that are different on the master and slave servers, or names that are too long to be legal column identifiers. A workaround is to supply aliases for columns that refer to local variables. Consider this statement when `myvar` has a value of 1:

```
CREATE TABLE t1 SELECT myvar;
```

That will be rewritten as follows:

```
CREATE TABLE t1 SELECT NAME_CONST(myvar, 1);
```

To ensure that the master and slave tables have the same column names, write the statement like this:

```
CREATE TABLE t1 SELECT myvar AS myvar;
```

The rewritten statement becomes:

```
CREATE TABLE t1 SELECT NAME_CONST(myvar, 1) AS myvar;
```

- A statement to be logged might contain references to user-defined variables. To handle this, MySQL writes a `SET` statement to the binary log to make sure that the variable exists on the slave with the same value as on the master. For example, if a statement refers to a variable `@my_var`, that statement will be preceded in the binary log by the following statement, where *`value`* is the value of `@my_var` on the master:

  ```
  SET @my_var = value;
  ```

- Procedure calls can occur within a committed or rolled-back transaction. Transactional context is accounted for so that the transactional aspects of procedure execution are replicated correctly. That is, the server logs those statements within the procedure that actually execute and modify data, and also logs `BEGIN`, `COMMIT`, and `ROLLBACK` statements as necessary. For example, if a procedure updates only transactional tables and is executed within a transaction that is rolled back, those updates are not logged. If the procedure occurs within a committed transaction, `BEGIN` and `COMMIT` statements are logged with the updates. For a procedure that executes within a rolled-back transaction, its statements are logged using the same rules that would apply if the statements were executed in standalone fashion:

  - Updates to transactional tables are not logged.
  - Updates to nontransactional tables are logged because rollback does not cancel them.
  - Updates to a mix of transactional and nontransactional tables are logged surrounded by `BEGIN` and `ROLLBACK` so that slaves will make the same changes and rollbacks as on the master.

- A stored procedure call is *not* written to the binary log at the statement level if the procedure is invoked from within a stored function. In that case, the only thing logged is the statement that invokes the function (if it occurs within a statement that is logged) or a `DO` statement (if it occurs within a statement that is not logged). For this reason, care should be exercised in the use of stored functions that invoke a procedure, even if the procedure is otherwise safe in itself.

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`INFORMATION_SCHEMA` provides access to database *metadata*, information about the MySQL server such as the name of a database or table, the data type of a column, or access privileges. Other terms that are sometimes used for this information are *data dictionary* and *system catalog*.

INFORMATION_SCHEMA Usage Notes

`INFORMATION_SCHEMA` is a database within each MySQL instance, the place that stores information about all the other databases that the MySQL server maintains. The `INFORMATION_SCHEMA` database contains several read-only tables. They are actually views, not base tables, so there are no files associated with them, and you cannot set triggers on them. Also, there is no database directory with that name.

Although you can select `INFORMATION_SCHEMA` as the default database with a `USE` statement, you can only read the contents of tables, not perform `INSERT`, `UPDATE`, or `DELETE` operations on them.

Example

Here is an example of a statement that retrieves information from `INFORMATION_SCHEMA`:

```
mysql> SELECT table_name, table_type, engine
       FROM information_schema.tables
       WHERE table_schema = 'db5'
       ORDER BY table_name;
+------------+------------+--------+
| table_name | table_type | engine |
+------------+------------+--------+
| fk         | BASE TABLE | InnoDB |
| fk2        | BASE TABLE | InnoDB |
| goto       | BASE TABLE | MyISAM |
| into       | BASE TABLE | MyISAM |
| k          | BASE TABLE | MyISAM |
| kurs       | BASE TABLE | MyISAM |
| loop       | BASE TABLE | MyISAM |
| pk         | BASE TABLE | InnoDB |
| t          | BASE TABLE | MyISAM |
| t2         | BASE TABLE | MyISAM |
| t3         | BASE TABLE | MyISAM |
| t7         | BASE TABLE | MyISAM |
| tables     | BASE TABLE | MyISAM |
| v          | VIEW       | NULL   |
| v2         | VIEW       | NULL   |
| v3         | VIEW       | NULL   |
| v56        | VIEW       | NULL   |
+------------+------------+--------+
17 rows in set (0.01 sec)
```

Explanation: The statement requests a list of all the tables in database `db5`, showing just three pieces of information: the name of the table, its type, and its storage engine.

Character Set Considerations

The definition for character columns (for example, `TABLES.TABLE_NAME`) is generally `VARCHAR(N) CHARACTER SET utf8` where *N* is at least 64. MySQL uses the default collation for this character set (`utf8_general_ci`) for all searches, sorts, comparisons, and other string operations on such columns.

Because some MySQL objects are represented as files, searches in `INFORMATION_SCHEMA` string columns can be affected by file system case sensitivity. For more information, see Section 10.1.8.7, “Using Collation in INFORMATION_SCHEMA Searches”.

INFORMATION_SCHEMA as Alternative to SHOW Statements

The `SELECT ... FROM INFORMATION_SCHEMA` statement is intended as a more consistent way to provide access to the information provided by the various `SHOW` statements that MySQL supports (`SHOW DATABASES`, `SHOW TABLES`, and so forth). Using `SELECT` has these advantages, compared to `SHOW`:

- It conforms to Codd's rules, because all access is done on tables.
- You can use the familiar syntax of the `SELECT` statement, and only need to learn some table and column names.
- The implementor need not worry about adding keywords.
- You can filter, sort, concatenate, and transform the results from `INFORMATION_SCHEMA` queries into whatever format your application needs, such as a data structure or a text representation to parse.
- This technique is more interoperable with other database systems. For example, Oracle Database users are familiar with querying tables in the Oracle data dictionary.

Because `SHOW` is familiar and widely used, the `SHOW` statements remain as an alternative. In fact, along with the implementation of `INFORMATION_SCHEMA`, there are enhancements to `SHOW` as described in Section 24.33, “Extensions to SHOW Statements”.

Privileges

Each MySQL user has the right to access these tables, but can see only the rows in the tables that correspond to objects for which the user has the proper access privileges. In some cases (for example, the `ROUTINE_DEFINITION` column in the `INFORMATION_SCHEMA.ROUTINES` table), users who have insufficient privileges see `NULL`. These restrictions do not apply for `InnoDB` tables; you can see them with only the `PROCESS` privilege.

The same privileges apply to selecting information from `INFORMATION_SCHEMA` and viewing the same information through `SHOW` statements. In either case, you must have some privilege on an object to see information about it.

Performance Considerations

`INFORMATION_SCHEMA` queries that search for information from more than one database might take a long time and impact performance. To check the efficiency of a query, you can use `EXPLAIN`. For information about using `EXPLAIN` output to tune `INFORMATION_SCHEMA` queries, see Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”.

Standards Considerations

The implementation for the `INFORMATION_SCHEMA` table structures in MySQL follows the ANSI/ISO SQL:2003 standard Part 11 *Schemata*. Our intent is approximate compliance with SQL:2003 core feature F021 *Basic information schema*.

Users of SQL Server 2000 (which also follows the standard) may notice a strong similarity. However, MySQL has omitted many columns that are not relevant for our implementation, and added columns that are MySQL-specific. One such column is the `ENGINE` column in the `INFORMATION_SCHEMA.TABLES` table.

Although other DBMSs use a variety of names, like `syscat` or `system`, the standard name is `INFORMATION_SCHEMA`.

To avoid using any name that is reserved in the standard or in DB2, SQL Server, or Oracle, we changed the names of some columns marked “MySQL extension”. (For example, we changed `COLLATION` to `TABLE_COLLATION` in the `TABLES` table.) See the list of reserved words near the end of this article: https://web.archive.org/web/20070428032454/http://www.dbazine.com/db2/db2-disarticles/gulutzan5.

Conventions in the INFORMATION_SCHEMA Reference Sections

The following sections describe each of the tables and columns in `INFORMATION_SCHEMA`. For each column, there are three pieces of information:

- “`INFORMATION_SCHEMA` Name” indicates the name for the column in the `INFORMATION_SCHEMA` table. This corresponds to the standard SQL name unless the “Remarks” field says “MySQL extension.”
- “`SHOW` Name” indicates the equivalent field name in the closest `SHOW` statement, if there is one.
- “Remarks” provides additional information where applicable. If this field is `NULL`, it means that the value of the column is always `NULL`. If this field says “MySQL extension,” the column is a MySQL extension to standard SQL.

Many sections indicate what `SHOW` statement is equivalent to a `SELECT` that retrieves information from `INFORMATION_SCHEMA`. For `SHOW` statements that display information for the default database if you omit a `FROM` *`db_name`* clause, you can often select information for the default database by adding an `AND TABLE_SCHEMA = SCHEMA()` condition to the `WHERE` clause of a query that retrieves information from an `INFORMATION_SCHEMA` table.

Related Information

These sections discuss additional `INFORMATION_SCHEMA`-related topics:

- information about `INFORMATION_SCHEMA` tables specific to the `InnoDB` storage engine: Section 24.31, “InnoDB INFORMATION_SCHEMA Tables”
- Answers to questions that are often asked concerning the `INFORMATION_SCHEMA` database: Section A.7, “MySQL 8.0 FAQ: INFORMATION_SCHEMA”
- `INFORMATION_SCHEMA` queries and the optimizer: Section 8.2.3, “Optimizing INFORMATION_SCHEMA Queries”
- The effect of collation on `INFORMATION_SCHEMA` comparisons: Section 10.1.8.7, “Using Collation in INFORMATION_SCHEMA Searches”

24.1 The INFORMATION_SCHEMA CHARACTER_SETS Table

The `CHARACTER_SETS` table provides information about available character sets.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`CHARACTER_SET_NAME`	`Charset`	
`DEFAULT_COLLATE_NAME`	`Default collation`	
`DESCRIPTION`	`Description`	MySQL extension
`MAXLEN`	`Maxlen`	MySQL extension

The following statements are equivalent:

```
SELECT * FROM INFORMATION_SCHEMA.CHARACTER_SETS
  [WHERE CHARACTER_SET_NAME LIKE 'wild']

SHOW CHARACTER SET
  [LIKE 'wild']
```

24.2 The INFORMATION_SCHEMA COLLATIONS Table

The `COLLATIONS` table provides information about collations for each character set.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`COLLATION_NAME`	`Collation`	
`CHARACTER_SET_NAME`	`Charset`	MySQL extension
`ID`	`Id`	MySQL extension
`IS_DEFAULT`	`Default`	MySQL extension
`IS_COMPILED`	`Compiled`	MySQL extension
`SORTLEN`	`Sortlen`	MySQL extension
`PAD_ATTRIBUTE`		omit

- `COLLATION_NAME` is the collation name.
- `CHARACTER_SET_NAME` is the name of the character set with which the collation is associated.
- `ID` is the collation ID.
- `IS_DEFAULT` indicates whether the collation is the default for its character set.
- `IS_COMPILED` indicates whether the character set is compiled into the server.
- `SORTLEN` is related to the amount of memory required to sort strings expressed in the character set.
- `PAD_ATTRIBUTE` indicates the collation pad attribute.

Collation information is also available from the `SHOW COLLATION` statement. The following statements are equivalent:

```
SELECT COLLATION_NAME FROM INFORMATION_SCHEMA.COLLATIONS
  [WHERE COLLATION_NAME LIKE 'wild']

SHOW COLLATION
  [LIKE 'wild']
```

24.3 The INFORMATION_SCHEMA COLLATION_CHARACTER_SET_APPLICABILITY Table

The `COLLATION_CHARACTER_SET_APPLICABILITY` table indicates what character set is applicable for what collation. The columns are equivalent to the first two display fields that we get from `SHOW COLLATION`.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`COLLATION_NAME`	`Collation`	
`CHARACTER_SET_NAME`	`Charset`	

24.4 The INFORMATION_SCHEMA COLUMNS Table

The COLUMNS table provides information about columns in tables.

The related ST_GEOMETRY_COLUMNS table is a view on COLUMNS that provides information about table columns that store spatial data. See Section 24.22, "The INFORMATION_SCHEMA ST_GEOMETRY_COLUMNS Table".

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLE_CATALOG		def
TABLE_SCHEMA		
TABLE_NAME		
COLUMN_NAME	Field	
ORDINAL_POSITION		see notes
COLUMN_DEFAULT	Default	
IS_NULLABLE	Null	
DATA_TYPE	Type	
CHARACTER_MAXIMUM_LENGTH	Type	
CHARACTER_OCTET_LENGTH		
NUMERIC_PRECISION	Type	
NUMERIC_SCALE	Type	
DATETIME_PRECISION	Type	
CHARACTER_SET_NAME		
COLLATION_NAME	Collation	
COLUMN_TYPE	Type	MySQL extension
COLUMN_KEY	Key	MySQL extension
EXTRA	Extra	MySQL extension
PRIVILEGES	Privileges	MySQL extension
COLUMN_COMMENT	Comment	MySQL extension
GENERATION_EXPRESSION		MySQL extension

Notes:

- In SHOW, the Type display includes values from several different COLUMNS columns.
- ORDINAL_POSITION is necessary because you might want to say ORDER BY ORDINAL_POSITION. Unlike SHOW, SELECT does not have automatic ordering.
- CHARACTER_OCTET_LENGTH should be the same as CHARACTER_MAXIMUM_LENGTH, except for multibyte character sets.
- CHARACTER_SET_NAME can be derived from Collation. For example, if you say SHOW FULL COLUMNS FROM t, and you see in the Collation column a value of utf8_swedish_ci, the character set is what is before the first underscore: utf8.
- GENERATION_EXPRESSION is nonempty for generated columns and displays the expression used to compute column values. For information about generated columns, see Section 13.1.15.8, "CREATE TABLE and Generated Columns".
- The EXTRA column contains VIRTUAL GENERATED or VIRTUAL STORED for generated columns.

The following statements are nearly equivalent:

```
SELECT COLUMN_NAME, DATA_TYPE, IS_NULLABLE, COLUMN_DEFAULT
```

```
  FROM INFORMATION_SCHEMA.COLUMNS
  WHERE table_name = 'tbl_name'
  [AND table_schema = 'db_name']
  [AND column_name LIKE 'wild']

SHOW COLUMNS
  FROM tbl_name
  [FROM db_name]
  [LIKE 'wild']
```

24.5 The INFORMATION_SCHEMA COLUMN_PRIVILEGES Table

The `COLUMN_PRIVILEGES` table provides information about column privileges. This information comes from the `mysql.columns_priv` grant table.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`GRANTEE`		`'user_name'@'host_name'` value
`TABLE_CATALOG`		`def`
`TABLE_SCHEMA`		
`TABLE_NAME`		
`COLUMN_NAME`		
`PRIVILEGE_TYPE`		
`IS_GRANTABLE`		

Notes:

- In the output from `SHOW FULL COLUMNS`, the privileges are all in one field and in lowercase, for example, `select,insert,update,references`. In `COLUMN_PRIVILEGES`, there is one privilege per row, in uppercase.
- `PRIVILEGE_TYPE` can contain one (and only one) of these values: `SELECT`, `INSERT`, `UPDATE`, `REFERENCES`.
- If the user has `GRANT OPTION` privilege, `IS_GRANTABLE` should be `YES`. Otherwise, `IS_GRANTABLE` should be `NO`. The output does not list `GRANT OPTION` as a separate privilege.

The following statements are *not* equivalent:

```
SELECT ... FROM INFORMATION_SCHEMA.COLUMN_PRIVILEGES

SHOW GRANTS ...
```

24.6 The INFORMATION_SCHEMA COLUMN_STATISTICS Table

The `COLUMN_STATISTICS` table is a view that provides access to information in the `column_statistics` data dictionary table, which stores histogram statistics about column values. It was added in MySQL 8.0.2.

For information about histogram statistics, see Section 8.9.6, “Optimizer Statistics”, and Section 13.7.2.1, “ANALYZE TABLE Syntax”.

You can see information only for columns for which you have some privilege.

`COLUMN_STATISTICS` has these columns:

- `SCHEMA_NAME`

  The names of the schema for which the statistics apply.

- `TABLE_NAME`

  The names of the column for which the statistics apply.

- `COLUMN_NAME`

  The names of the column for which the statistics apply.

- `HISTOGRAM`

  A `JSON` object describing the column statistics, stored as a histogram.

24.7 The INFORMATION_SCHEMA ENGINES Table

The `ENGINES` table provides information about storage engines.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`ENGINE`	`Engine`	MySQL extension
`SUPPORT`	`Support`	MySQL extension
`COMMENT`	`Comment`	MySQL extension
`TRANSACTIONS`	`Transactions`	MySQL extension
`XA`	`XA`	MySQL extension
`SAVEPOINTS`	`Savepoints`	MySQL extension

Notes:

- The `ENGINES` table is a nonstandard table. Its contents correspond to the columns of the `SHOW ENGINES` statement. For descriptions of its columns, see Section 13.7.5.16, “SHOW ENGINES Syntax”.

See also Section 13.7.5.16, “SHOW ENGINES Syntax”.

24.8 The INFORMATION_SCHEMA EVENTS Table

The `EVENTS` table provides information about scheduled events, which are discussed in Section 23.4, “Using the Event Scheduler”. The `SHOW Name` values correspond to column names of the `SHOW EVENTS` statement.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`EVENT_CATALOG`		`def`, MySQL extension
`EVENT_SCHEMA`	`Db`	MySQL extension
`EVENT_NAME`	`Name`	MySQL extension
`DEFINER`	`Definer`	MySQL extension
`TIME_ZONE`	`Time zone`	MySQL extension
`EVENT_BODY`		MySQL extension
`EVENT_DEFINITION`		MySQL extension
`EVENT_TYPE`	`Type`	MySQL extension
`EXECUTE_AT`	`Execute at`	MySQL extension
`INTERVAL_VALUE`	`Interval value`	MySQL extension
`INTERVAL_FIELD`	`Interval field`	MySQL extension
`SQL_MODE`		MySQL extension
`STARTS`	`Starts`	MySQL extension

INFORMATION_SCHEMA Name	SHOW Name	Remarks
ENDS	Ends	MySQL extension
STATUS	Status	MySQL extension
ON_COMPLETION		MySQL extension
CREATED		MySQL extension
LAST_ALTERED		MySQL extension
LAST_EXECUTED		MySQL extension
EVENT_COMMENT		MySQL extension
ORIGINATOR	Originator	MySQL extension
CHARACTER_SET_CLIENT	character_set_client	MySQL extension
COLLATION_CONNECTION	collation_connection	MySQL extension
DATABASE_COLLATION	Database Collation	MySQL extension

Notes:

- The EVENTS table is a nonstandard table.
- EVENT_CATALOG: The value of this column is always def.
- EVENT_SCHEMA: The name of the schema (database) to which this event belongs.
- EVENT_NAME: The name of the event.
- DEFINER: The account of the user who created the event, in *'user_name'@'host_name'* format.
- TIME_ZONE: The event time zone, which is the time zone used for scheduling the event and that is in effect within the event as it executes. The default value is SYSTEM.
- EVENT_BODY: The language used for the statements in the event's DO clause; in MySQL 8.0, this is always SQL.

  This column is not to be confused with the column of the same name (now named EVENT_DEFINITION) that existed in earlier MySQL versions.
- EVENT_DEFINITION: The text of the SQL statement making up the event's DO clause; in other words, the statement executed by this event.
- EVENT_TYPE: The event repetition type, either ONE TIME (transient) or RECURRING (repeating).
- EXECUTE_AT: For a one-time event, this is the DATETIME value specified in the AT clause of the CREATE EVENT statement used to create the event, or of the last ALTER EVENT statement that modified the event. The value shown in this column reflects the addition or subtraction of any INTERVAL value included in the event's AT clause. For example, if an event is created using ON SCHEDULE AT CURRENT_TIMESTAMP + '1:6' DAY_HOUR, and the event was created at 2006-02-09 14:05:30, the value shown in this column would be '2006-02-10 20:05:30'.

  If the event's timing is determined by an EVERY clause instead of an AT clause (that is, if the event is recurring), the value of this column is NULL.
- INTERVAL_VALUE: For recurring events, this column contains the numeric portion of the event's EVERY clause.

  For a one-time event (that is, an event whose timing is determined by an AT clause), this column is NULL.
- INTERVAL_FIELD: For recurring events, this column contains the units portion of the EVERY clause governing the timing of the event. Thus, this column contains a value such as 'YEAR', 'QUARTER', 'DAY', and so on.

For a one-time event (that is, an event whose timing is determined by an `AT` clause), this column is `NULL`.

- `SQL_MODE`: The SQL mode in effect when the event was created or altered, and under which the event executes. For the permitted values, see Section 5.1.8, “Server SQL Modes”.
- `STARTS`: For a recurring event whose definition includes a `STARTS` clause, this column contains the corresponding `DATETIME` value. As with the `EXECUTE_AT` column, this value resolves any expressions used.

  If there is no `STARTS` clause affecting the timing of the event, this column is `NULL`
- `ENDS`: For a recurring event whose definition includes a `ENDS` clause, this column contains the corresponding `DATETIME` value. As with the `EXECUTE_AT` column, this value resolves any expressions used.

  If there is no `ENDS` clause affecting the timing of the event, this column is `NULL`.
- `STATUS`: One of the three values `ENABLED`, `DISABLED`, or `SLAVESIDE_DISABLED`.

  `SLAVESIDE_DISABLED` indicates that the creation of the event occurred on another MySQL server acting as a replication master and was replicated to the current MySQL server which is acting as a slave, but the event is not presently being executed on the slave. See Section 18.4.1.12, “Replication of Invoked Features”, for more information.
- `ON_COMPLETION`: One of the two values `PRESERVE` or `NOT PRESERVE`.
- `CREATED`: The date and time when the event was created. This is a `TIMESTAMP` value.
- `LAST_ALTERED`: The date and time when the event was last modified. This is a `TIMESTAMP` value. If the event has not been modified since its creation, this column holds the same value as the `CREATED` column.
- `LAST_EXECUTED`: The date and time when the event last executed. A `DATETIME` value. If the event has never executed, this column is `NULL`.

  `LAST_EXECUTED` indicates when the event started. As a result, the `ENDS` column is never less than `LAST_EXECUTED`.
- `EVENT_COMMENT`: The text of a comment, if the event has one. If not, the value of this column is an empty string.
- `ORIGINATOR`: The server ID of the MySQL server on which the event was created; used in replication. The default value is 0.
- `CHARACTER_SET_CLIENT`: The session value of the `character_set_client` system variable when the event was created.
- `COLLATION_CONNECTION`: The session value of the `collation_connection` system variable when the event was created.
- `DATABASE_COLLATION`: The collation of the database with which the event is associated.

Example: Suppose that the user `jon@ghidora` creates an event named `e_daily`, and then modifies it a few minutes later using an `ALTER EVENT` statement, as shown here:

```
DELIMITER |

CREATE EVENT e_daily
    ON SCHEDULE
      EVERY 1 DAY
    COMMENT 'Saves total number of sessions then clears the table each day'
```

```
    DO
      BEGIN
        INSERT INTO site_activity.totals (time, total)
          SELECT CURRENT_TIMESTAMP, COUNT(*)
            FROM site_activity.sessions;
        DELETE FROM site_activity.sessions;
      END |

DELIMITER ;

ALTER EVENT e_daily
    ENABLE;
```

(Note that comments can span multiple lines.)

This user can then run the following `SELECT` statement, and obtain the output shown:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.EVENTS
     > WHERE EVENT_NAME = 'e_daily'
     > AND EVENT_SCHEMA = 'myschema'\G
*************************** 1. row ***************************
       EVENT_CATALOG: def
        EVENT_SCHEMA: test
          EVENT_NAME: e_daily
             DEFINER: me@localhost
           TIME_ZONE: SYSTEM
          EVENT_BODY: SQL
    EVENT_DEFINITION: BEGIN
        INSERT INTO site_activity.totals (time, total)
          SELECT CURRENT_TIMESTAMP, COUNT(*)
            FROM site_activity.sessions;
        DELETE FROM site_activity.sessions;
      END
          EVENT_TYPE: RECURRING
          EXECUTE_AT: NULL
      INTERVAL_VALUE: 1
      INTERVAL_FIELD: DAY
            SQL_MODE:
              STARTS: 2008-09-03 12:13:39
                ENDS: NULL
              STATUS: ENABLED
       ON_COMPLETION: NOT PRESERVE
             CREATED: 2008-09-03 12:13:39
        LAST_ALTERED: 2008-09-03 12:13:39
       LAST_EXECUTED: NULL
       EVENT_COMMENT: Saves total number of sessions then clears the
                      table each day
          ORIGINATOR: 1
CHARACTER_SET_CLIENT: utf8
COLLATION_CONNECTION: utf8_swedish_ci
  DATABASE_COLLATION: utf8_swedish_ci
```

Times in the `EVENTS` table are displayed using the event time zone or the current session time zone, as described in Section 23.4.4, “Event Metadata”.

See also Section 13.7.5.18, “SHOW EVENTS Syntax”.

24.9 The INFORMATION_SCHEMA FILES Table

The `FILES` table provides information about the files in which MySQL tablespace data is stored.

`INFORMATION_SCHEMA.FILES` reports data about `NDB` and `InnoDB` data files.

`INFORMATION_SCHEMA` Name	**`SHOW` Name**	**Remarks**
`FILE_ID`		MySQL extension
`FILE_NAME`		MySQL extension
`FILE_TYPE`		MySQL extension

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLESPACE_NAME		MySQL extension
TABLE_CATALOG		MySQL extension
TABLE_SCHEMA		MySQL extension
TABLE_NAME		MySQL extension
LOGFILE_GROUP_NAME		MySQL extension
LOGFILE_GROUP_NUMBER		MySQL extension
ENGINE		MySQL extension
FULLTEXT_KEYS		MySQL extension
DELETED_ROWS		MySQL extension
UPDATE_COUNT		MySQL extension
FREE_EXTENTS		MySQL extension
TOTAL_EXTENTS		MySQL extension
EXTENT_SIZE		MySQL extension
INITIAL_SIZE		MySQL extension
MAXIMUM_SIZE		MySQL extension
AUTOEXTEND_SIZE		MySQL extension
CREATION_TIME		MySQL extension
LAST_UPDATE_TIME		MySQL extension
LAST_ACCESS_TIME		MySQL extension
RECOVER_TIME		MySQL extension
TRANSACTION_COUNTER		MySQL extension
VERSION		MySQL extension
ROW_FORMAT		MySQL extension
TABLE_ROWS		MySQL extension
AVG_ROW_LENGTH		MySQL extension
DATA_LENGTH		MySQL extension
MAX_DATA_LENGTH		MySQL extension
INDEX_LENGTH		MySQL extension
DATA_FREE		MySQL extension
CREATE_TIME		MySQL extension
UPDATE_TIME		MySQL extension
CHECK_TIME		MySQL extension
CHECKSUM		MySQL extension
STATUS		MySQL extension
EXTRA		MySQL extension

InnoDB Notes:

The following notes apply to `InnoDB` data files. `INFORMATION_SCHEMA.FILES` fields that are not described below are not applicable to `InnoDB` and report a `NULL` value.

- Data reported by `INFORMATION_SCHEMA.FILES` is reported from the `InnoDB` in-memory cache for open files. By comparison, `INFORMATION_SCHEMA.INNODB_SYS_DATAFILES` reports data from the `InnoDB` `SYS_DATAFILES` internal data dictionary table.

- The data reported by `INFORMATION_SCHEMA.FILES` includes temporary tablespace data. This data is not available in the internal `SYS_DATAFILES` data dictionary table, and is therefore not reported by `INNODB_SYS_DATAFILES`.
- Undo tablespace data is reported by `INFORMATION_SCHEMA.FILES` when separate undo tablespaces are configured using the `innodb_undo_tablespaces` configuration option.
- `FILE_ID` is the tablespace ID, also referred to as the `space_id` or `fil_space_t::id`.
- `FILE_NAME` is the name of the data file. File-per-table and general tablespaces have a `.ibd` file name extension. Undo tablespaces are prefixed by `undo`. The system tablespace is prefixed by `ibdata`. Temporary tablespaces are prefixed by `ibtmp`. The file name includes the file path, which may be relative to the MySQL data directory (`datadir`).
- `FILE_TYPE` is the tablespace file type. There are three possible file types for `InnoDB` files. `TABLESPACE` is the file type for any system, general, or file-per-table tablespace file that holds tables, indexes, or other forms of user data. `TEMPORARY` is the file type for temporary tablespaces. `UNDO LOG` is the file type for undo tablespaces, which hold undo records. Separate undo tablespaces are configured using the `innodb_undo_tablespaces` option.
- `TABLESPACE_NAME` is the SQL name for the tablespace. A general tablespace name is the `SYS_TABLESPACES.NAME` value. For other tablespace files, names start with `innodb_`, such as `innodb_system`, `innodb_undo`, and `innodb_file_per_table`. The file-per-table tablespace name format is `innodb_file_per_table_##`, where `##` is the tablespace ID.
- `ENGINE` is the storage engine. For `InnoDB` files, the value is always `InnoDB`.
- `FREE_EXTENTS` is the number of fully free extents in the current data file.
- `TOTAL_EXTENTS` is the number of full extents used in the current data file. Any partial extent at the end of the file is not counted.
- `EXTENT_SIZE` is 1048576 (1MB) for files with a 4k, 8k, or 16k page size. Extent size is 2097152 bytes (2MB) for files with a 32k page size, and 4194304 (4MB) for files with a 64k page size. `INFORMATION_SCHEMA.FILES` does not report `InnoDB` page size. Page size is defined by the `innodb_page_size` option. Extent size information can also be retrieved from `INNODB_SYS_TABLESPACES` where `FILES.FILE_ID = INNODB_SYS_TABLESPACES.SPACE_ID`.
- `INITIAL_SIZE` is the initial size of the file, in bytes.
- `MAXIMUM_SIZE` is the maximum number of bytes allowed in the file. The value is `NULL` for all data files except for predefined system tablespace data files. Maximum system tablespace file size is defined by `innodb_data_file_path`. Maximum temporary tablespace file size is defined by `innodb_temp_data_file_path`. A `NULL` value for a predefined system tablespace data file indicates that a file size limit was not defined explicitly.
- `AUTOEXTEND_SIZE` is the auto-extend size defined by `innodb_data_file_path` for the system tablespace, or defined by `innodb_temp_data_file_path` for temporary tablespaces.
- `DATA_FREE` is the total amount of free space (in bytes) for the entire tablespace. Predefined system tablespaces, which include the system tablespace and temporary table tablespaces, may have one or more data files.
- `STATUS` is `NORMAL` by default. `InnoDB` file-per-table tablespaces may report `IMPORTING`, which indicates that the tablespace is not yet available.
- The following query returns all data pertinent to `InnoDB` tablespaces.

```
mysql> SELECT FILE_ID, FILE_NAME, FILE_TYPE, TABLESPACE_NAME, FREE_EXTENTS, TOTAL_EXTENTS,
EXTENT_SIZE, INITIAL_SIZE, MAXIMUM_SIZE, AUTOEXTEND_SIZE, DATA_FREE, STATUS ENGINE
FROM INFORMATION_SCHEMA.FILES \G
```

24.10 The INFORMATION_SCHEMA KEY_COLUMN_USAGE Table

The KEY_COLUMN_USAGE table describes which key columns have constraints.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
CONSTRAINT_CATALOG		def
CONSTRAINT_SCHEMA		
CONSTRAINT_NAME		
TABLE_CATALOG		def
TABLE_SCHEMA		
TABLE_NAME		
COLUMN_NAME		
ORDINAL_POSITION		
POSITION_IN_UNIQUE_CONSTRAINT		
REFERENCED_TABLE_SCHEMA		
REFERENCED_TABLE_NAME		
REFERENCED_COLUMN_NAME		

Notes:

- If the constraint is a foreign key, then this is the column of the foreign key, not the column that the foreign key references.
- The value of ORDINAL_POSITION is the column's position within the constraint, not the column's position within the table. Column positions are numbered beginning with 1.
- The value of POSITION_IN_UNIQUE_CONSTRAINT is NULL for unique and primary-key constraints. For foreign-key constraints, it is the ordinal position in key of the table that is being referenced.

  Suppose that there are two tables name t1 and t3 that have the following definitions:

```
CREATE TABLE t1
(
    s1 INT,
    s2 INT,
    s3 INT,
    PRIMARY KEY(s3)
) ENGINE=InnoDB;

CREATE TABLE t3
(
    s1 INT,
    s2 INT,
    s3 INT,
    KEY(s1),
    CONSTRAINT CO FOREIGN KEY (s2) REFERENCES t1(s3)
) ENGINE=InnoDB;
```

  For those two tables, the KEY_COLUMN_USAGE table has two rows:

  - One row with CONSTRAINT_NAME = 'PRIMARY', TABLE_NAME = 't1', COLUMN_NAME = 's3', ORDINAL_POSITION = 1, POSITION_IN_UNIQUE_CONSTRAINT = NULL.
  - One row with CONSTRAINT_NAME = 'CO', TABLE_NAME = 't3', COLUMN_NAME = 's2', ORDINAL_POSITION = 1, POSITION_IN_UNIQUE_CONSTRAINT = 1.

24.11 The INFORMATION_SCHEMA OPTIMIZER_TRACE Table

The OPTIMIZER_TRACE table provides information produced by the optimizer tracing capability. To enable tracking, use the optimizer_trace system variable. For details, see MySQL Internals: Tracing the Optimizer.

24.12 The INFORMATION_SCHEMA PARAMETERS Table

The PARAMETERS table provides information about stored procedure and function parameters, and about return values for stored functions.

INFORMATION_SCHEMA Name	Remarks
SPECIFIC_CATALOG	def
SPECIFIC_SCHEMA	routine database
SPECIFIC_NAME	routine name
ORDINAL_POSITION	1, 2, 3, ... for parameters, 0 for function RETURNS clause
PARAMETER_MODE	IN, OUT, INOUT (NULL for RETURNS)
PARAMETER_NAME	parameter name (NULL for RETURNS)
DATA_TYPE	same as for COLUMNS table
CHARACTER_MAXIMUM_LENGTH	same as for COLUMNS table
CHARACTER_OCTET_LENGTH	same as for COLUMNS table
NUMERIC_PRECISION	same as for COLUMNS table
NUMERIC_SCALE	same as for COLUMNS table
DATETIME_PRECISION	same as for COLUMNS table
CHARACTER_SET_NAME	same as for COLUMNS table
COLLATION_NAME	same as for COLUMNS table
DTD_IDENTIFIER	same as for COLUMNS table
ROUTINE_TYPE	same as for ROUTINES table

Notes:

- For successive parameters of a stored procedure or function, the ORDINAL_POSITION values are 1, 2, 3, and so forth. For a stored function, there is also a row that describes the data type for the RETURNS clause. The return value is not a true parameter, so the row that describes it has these unique characteristics:
  - The ORDINAL_POSITION value is 0.
  - The PARAMETER_NAME and PARAMETER_MODE values are NULL because the return value has no name and the mode does not apply.

24.13 The INFORMATION_SCHEMA PARTITIONS Table

The PARTITIONS table provides information about table partitions. See Chapter 22, *Partitioning*, for more information about partitioning tables.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLE_CATALOG		MySQL extension
TABLE_SCHEMA		MySQL extension
TABLE_NAME		MySQL extension
PARTITION_NAME		MySQL extension

INFORMATION_SCHEMA Name	SHOW Name	Remarks
SUBPARTITION_NAME		MySQL extension
PARTITION_ORDINAL_POSITION		MySQL extension
SUBPARTITION_ORDINAL_POSITION		MySQL extension
PARTITION_METHOD		MySQL extension
SUBPARTITION_METHOD		MySQL extension
PARTITION_EXPRESSION		MySQL extension
SUBPARTITION_EXPRESSION		MySQL extension
PARTITION_DESCRIPTION		MySQL extension
TABLE_ROWS		MySQL extension
AVG_ROW_LENGTH		MySQL extension
DATA_LENGTH		MySQL extension
MAX_DATA_LENGTH		MySQL extension
INDEX_LENGTH		MySQL extension
DATA_FREE		MySQL extension
CREATE_TIME		MySQL extension
UPDATE_TIME		MySQL extension
CHECK_TIME		MySQL extension
CHECKSUM		MySQL extension
PARTITION_COMMENT		MySQL extension
NODEGROUP		MySQL extension
TABLESPACE_NAME		MySQL extension

Notes:

- The PARTITIONS table is a nonstandard table.

  Each record in this table corresponds to an individual partition or subpartition of a partitioned table.

- TABLE_CATALOG: This column is always def.
- TABLE_SCHEMA: This column contains the name of the database to which the table belongs.
- TABLE_NAME: This column contains the name of the table containing the partition.
- PARTITION_NAME: The name of the partition.
- SUBPARTITION_NAME: If the PARTITIONS table record represents a subpartition, then this column contains the name of subpartition; otherwise it is NULL.
- PARTITION_ORDINAL_POSITION: All partitions are indexed in the same order as they are defined, with 1 being the number assigned to the first partition. The indexing can change as partitions are added, dropped, and reorganized; the number shown is this column reflects the current order, taking into account any indexing changes.
- SUBPARTITION_ORDINAL_POSITION: Subpartitions within a given partition are also indexed and reindexed in the same manner as partitions are indexed within a table.
- PARTITION_METHOD: One of the values RANGE, LIST, HASH, LINEAR HASH, KEY, or LINEAR KEY; that is, one of the available partitioning types as discussed in Section 22.2, “Partitioning Types”.
- SUBPARTITION_METHOD: One of the values HASH, LINEAR HASH, KEY, or LINEAR KEY; that is, one of the available subpartitioning types as discussed in Section 22.2.6, “Subpartitioning”.

- PARTITION_EXPRESSION: This is the expression for the partitioning function used in the CREATE TABLE or ALTER TABLE statement that created the table's current partitioning scheme.

  For example, consider a partitioned table created in the test database using this statement:

```
CREATE TABLE tp (
    c1 INT,
    c2 INT,
    c3 VARCHAR(25)
)
PARTITION BY HASH(c1 + c2)
PARTITIONS 4;
```

  The PARTITION_EXPRESSION column in a PARTITIONS table record for a partition from this table displays c1 + c2, as shown here:

```
mysql> SELECT DISTINCT PARTITION_EXPRESSION
     >     FROM INFORMATION_SCHEMA.PARTITIONS
     >     WHERE TABLE_NAME='tp' AND TABLE_SCHEMA='test';
+----------------------+
| PARTITION_EXPRESSION |
+----------------------+
| c1 + c2              |
+----------------------+
1 row in set (0.09 sec)
```

- SUBPARTITION_EXPRESSION: This works in the same fashion for the subpartitioning expression that defines the subpartitioning for a table as PARTITION_EXPRESSION does for the partitioning expression used to define a table's partitioning.

  If the table has no subpartitions, then this column is NULL.

- PARTITION_DESCRIPTION: This column is used for RANGE and LIST partitions. For a RANGE partition, it contains the value set in the partition's VALUES LESS THAN clause, which can be either an integer or MAXVALUE. For a LIST partition, this column contains the values defined in the partition's VALUES IN clause, which is a comma-separated list of integer values.

  For partitions whose PARTITION_METHOD is other than RANGE or LIST, this column is always NULL.

- TABLE_ROWS: The number of table rows in the partition.

  For partitioned InnoDB tables, the row count given in the TABLE_ROWS column is only an estimated value used in SQL optimization, and may not always be exact.

- AVG_ROW_LENGTH: The average length of the rows stored in this partition or subpartition, in bytes.

  This is the same as DATA_LENGTH divided by TABLE_ROWS.

- DATA_LENGTH: The total length of all rows stored in this partition or subpartition, in bytes—that is, the total number of bytes stored in the partition or subpartition.

- MAX_DATA_LENGTH: The maximum number of bytes that can be stored in this partition or subpartition.

- INDEX_LENGTH: The length of the index file for this partition or subpartition, in bytes.

- DATA_FREE: The number of bytes allocated to the partition or subpartition but not used.

- CREATE_TIME: The time of the partition's or subpartition's creation.

- UPDATE_TIME: The time that the partition or subpartition was last modified.

- CHECK_TIME: The last time that the table to which this partition or subpartition belongs was checked.

For partitioned InnoDB tables, this column is always NULL.

- CHECKSUM: The checksum value, if any; otherwise, this column is NULL.
- PARTITION_COMMENT: This column contains the text of any comment made for the partition.

  In MySQL 8.0, the maximum length for a partition comment is defined as 1024 characters, and the display width of the PARTITION_COMMENT column is also 1024, characters to match this limit (Bug #11748924, Bug #37728).

  The default value for this column is an empty string.
- NODEGROUP: This is the nodegroup to which the partition belongs. This is relevant only to MySQL Cluster tables; otherwise the value of this column is always 0.
- TABLESPACE_NAME: This column contains the name of the tablespace to which the partition belongs. The value of this column is always DEFAULT.
- A nonpartitioned table has one record in INFORMATION_SCHEMA.PARTITIONS; however, the values of the PARTITION_NAME, SUBPARTITION_NAME, PARTITION_ORDINAL_POSITION, SUBPARTITION_ORDINAL_POSITION, PARTITION_METHOD, SUBPARTITION_METHOD, PARTITION_EXPRESSION, SUBPARTITION_EXPRESSION, and PARTITION_DESCRIPTION columns are all NULL. (The PARTITION_COMMENT column in this case is blank.)

24.14 The INFORMATION_SCHEMA PLUGINS Table

The PLUGINS table provides information about server plugins.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
PLUGIN_NAME	Name	MySQL extension
PLUGIN_VERSION		MySQL extension
PLUGIN_STATUS	Status	MySQL extension
PLUGIN_TYPE	Type	MySQL extension
PLUGIN_TYPE_VERSION		MySQL extension
PLUGIN_LIBRARY	Library	MySQL extension
PLUGIN_LIBRARY_VERSION		MySQL extension
PLUGIN_AUTHOR		MySQL extension
PLUGIN_DESCRIPTION		MySQL extension
PLUGIN_LICENSE	License	MySQL extension
LOAD_OPTION		MySQL extension

Notes:

- The PLUGINS table is a nonstandard table.
- PLUGIN_NAME is the name used to refer to the plugin in statements such as INSTALL PLUGIN and UNINSTALL PLUGIN.
- PLUGIN_VERSION is the version from the plugin's general type descriptor.
- PLUGIN_STATUS indicates the plugin status, one of ACTIVE, INACTIVE, DISABLED, or DELETED.
- PLUGIN_TYPE indicates the type of plugin, such as STORAGE ENGINE, INFORMATION_SCHEMA, or AUTHENTICATION.
- PLUGIN_TYPE_VERSION is the version from the plugin's type-specific descriptor.

- PLUGIN_LIBRARY is the name of the plugin shared library file. This is the name used to refer to the plugin file in statements such as INSTALL PLUGIN and UNINSTALL PLUGIN. This file is located in the directory named by the plugin_dir system variable. If the library name is NULL, the plugin is compiled in and cannot be uninstalled with UNINSTALL PLUGIN.
- PLUGIN_LIBRARY_VERSION indicates the plugin API interface version.
- PLUGIN_AUTHOR names the plugin author.
- PLUGIN_DESCRIPTION provides a short description of the plugin.
- PLUGIN_LICENSE indicates how the plugin is licensed; for example, GPL.
- LOAD_OPTION indicates how the plugin was loaded. The value is OFF, ON, FORCE, or FORCE_PLUS_PERMANENT. See Section 5.6.2, “Installing and Uninstalling Plugins”.

For plugins installed with INSTALL PLUGIN, the PLUGIN_NAME and PLUGIN_LIBRARY values are also registered in the mysql.plugin table.

These statements are equivalent:

```
SELECT
  PLUGIN_NAME, PLUGIN_STATUS, PLUGIN_TYPE,
  PLUGIN_LIBRARY, PLUGIN_LICENSE
FROM INFORMATION_SCHEMA.PLUGINS;

SHOW PLUGINS;
```

For information about plugin data structures that form the basis of the information in the PLUGINS table, see Section 28.2, “The MySQL Plugin API”.

Plugin information is also available using the SHOW PLUGINS statement. See Section 13.7.5.25, “SHOW PLUGINS Syntax”.

24.15 The INFORMATION_SCHEMA PROCESSLIST Table

The PROCESSLIST table provides information about which threads are running.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
ID	Id	MySQL extension
USER	User	MySQL extension
HOST	Host	MySQL extension
DB	db	MySQL extension
COMMAND	Command	MySQL extension
TIME	Time	MySQL extension
STATE	State	MySQL extension
INFO	Info	MySQL extension

For an extensive description of the table columns, see Section 13.7.5.29, “SHOW PROCESSLIST Syntax”.

Notes:

- The PROCESSLIST table is a nonstandard table.
- Like the output from the corresponding SHOW statement, the PROCESSLIST table will only show information about your own threads, unless you have the PROCESS privilege, in which case you will see information about other threads, too. As an anonymous user, you cannot see any rows at all.

- If an SQL statement refers to `INFORMATION_SCHEMA.PROCESSLIST`, MySQL populates the entire table once, when statement execution begins, so there is read consistency during the statement. There is no read consistency for a multi-statement transaction, though.
- Process information is also available from the `performance_schema.threads` table. However, access to `threads` does not require a mutex and has minimal impact on server performance. `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` have negative performance consequences because they require a mutex. `threads` also shows information about background threads, which `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` do not. This means that `threads` can be used to monitor activity the other thread information sources cannot.

The following statements are equivalent:

```
SELECT * FROM INFORMATION_SCHEMA.PROCESSLIST

SHOW FULL PROCESSLIST
```

24.16 The INFORMATION_SCHEMA PROFILING Table

The `PROFILING` table provides statement profiling information. Its contents correspond to the information produced by the `SHOW PROFILES` and `SHOW PROFILE` statements (see Section 13.7.5.31, "SHOW PROFILES Syntax"). The table is empty unless the `profiling` session variable is set to 1.

Note

This table is deprecated and will be removed in a future MySQL release. Use the Performance Schema instead; see Chapter 25, *MySQL Performance Schema*.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`QUERY_ID`	`Query_ID`	
`SEQ`		
`STATE`	`Status`	
`DURATION`	`Duration`	
`CPU_USER`	`CPU_user`	
`CPU_SYSTEM`	`CPU_system`	
`CONTEXT_VOLUNTARY`	`Context_voluntary`	
`CONTEXT_INVOLUNTARY`	`Context_involuntary`	
`BLOCK_OPS_IN`	`Block_ops_in`	
`BLOCK_OPS_OUT`	`Block_ops_out`	
`MESSAGES_SENT`	`Messages_sent`	
`MESSAGES_RECEIVED`	`Messages_received`	
`PAGE_FAULTS_MAJOR`	`Page_faults_major`	
`PAGE_FAULTS_MINOR`	`Page_faults_minor`	
`SWAPS`	`Swaps`	
`SOURCE_FUNCTION`	`Source_function`	
`SOURCE_FILE`	`Source_file`	
`SOURCE_LINE`	`Source_line`	

Notes:

- `QUERY_ID` is a numeric statement identifier.

- `SEQ` is a sequence number indicating the display order for rows with the same `QUERY_ID` value.
- `STATE` is the profiling state to which the row measurements apply.
- `DURATION` indicates how long statement execution remained in the given state, in seconds.
- `CPU_USER` and `CPU_SYSTEM` indicate user and system CPU use, in seconds.
- `CONTEXT_VOLUNTARY` and `CONTEXT_INVOLUNTARY` indicate how many voluntary and involuntary context switches occurred.
- `BLOCK_OPS_IN` and `BLOCK_OPS_OUT` indicate the number of block input and output operations.
- `MESSAGES_SENT` and `MESSAGES_RECEIVED` indicate the number of communication messages sent and received.
- `PAGE_FAULTS_MAJOR` and `PAGE_FAULTS_MINOR` indicate the number of major and minor page faults.
- `SWAPS` indicates how many swaps occurred.
- `SOURCE_FUNCTION`, `SOURCE_FILE`, and `SOURCE_LINE` provide information indicating where in the source code the profiled state executes.

24.17 The INFORMATION_SCHEMA REFERENTIAL_CONSTRAINTS Table

The `REFERENTIAL_CONSTRAINTS` table provides information about foreign keys.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
CONSTRAINT_CATALOG		def
CONSTRAINT_SCHEMA		
CONSTRAINT_NAME		
UNIQUE_CONSTRAINT_CATALOG		def
UNIQUE_CONSTRAINT_SCHEMA		
UNIQUE_CONSTRAINT_NAME		
MATCH_OPTION		
UPDATE_RULE		
DELETE_RULE		
TABLE_NAME		
REFERENCED_TABLE_NAME		

Notes:

- `TABLE_NAME` has the same value as `TABLE_NAME` in `INFORMATION_SCHEMA.TABLE_CONSTRAINTS`.
- `CONSTRAINT_SCHEMA` and `CONSTRAINT_NAME` identify the foreign key.
- `UNIQUE_CONSTRAINT_SCHEMA`, `UNIQUE_CONSTRAINT_NAME`, and `REFERENCED_TABLE_NAME` identify the referenced key.
- The only valid value at this time for `MATCH_OPTION` is `NONE`.
- The possible values for `UPDATE_RULE` or `DELETE_RULE` are `CASCADE`, `SET NULL`, `SET DEFAULT`, `RESTRICT`, `NO ACTION`.

24.18 The INFORMATION_SCHEMA ROUTINES Table

The `ROUTINES` table provides information about stored routines (both procedures and functions). The `ROUTINES` table does not include user-defined functions (UDFs).

`INFORMATION_SCHEMA` Name	Remarks
`SPECIFIC_NAME`	
`ROUTINE_CATALOG`	`def`
`ROUTINE_SCHEMA`	
`ROUTINE_NAME`	
`ROUTINE_TYPE`	`{PROCEDURE|FUNCTION}`
`DATA_TYPE`	same as for `COLUMNS` table
`CHARACTER_MAXIMUM_LENGTH`	same as for `COLUMNS` table
`CHARACTER_OCTET_LENGTH`	same as for `COLUMNS` table
`NUMERIC_PRECISION`	same as for `COLUMNS` table
`NUMERIC_SCALE`	same as for `COLUMNS` table
`DATETIME_PRECISION`	same as for `COLUMNS` table
`CHARACTER_SET_NAME`	same as for `COLUMNS` table
`COLLATION_NAME`	same as for `COLUMNS` table
`DTD_IDENTIFIER`	data type descriptor
`ROUTINE_BODY`	`SQL`
`ROUTINE_DEFINITION`	
`EXTERNAL_NAME`	`NULL`
`EXTERNAL_LANGUAGE`	`NULL`
`PARAMETER_STYLE`	`SQL`
`IS_DETERMINISTIC`	
`SQL_DATA_ACCESS`	
`SQL_PATH`	`NULL`
`SECURITY_TYPE`	
`CREATED`	
`LAST_ALTERED`	
`SQL_MODE`	MySQL extension
`ROUTINE_COMMENT`	MySQL extension
`DEFINER`	MySQL extension
`CHARACTER_SET_CLIENT`	MySQL extension
`COLLATION_CONNECTION`	MySQL extension
`DATABASE_COLLATION`	MySQL extension

Notes:

- MySQL calculates `EXTERNAL_LANGUAGE` thus:
  - If `mysql.proc.language='SQL'`, `EXTERNAL_LANGUAGE` is `NULL`
  - Otherwise, `EXTERNAL_LANGUAGE` is what is in `mysql.proc.language`. However, we do not have external languages yet, so it is always `NULL`.

- `CREATED`: The date and time when the routine was created. This is a `TIMESTAMP` value.
- `LAST_ALTERED`: The date and time when the routine was last modified. This is a `TIMESTAMP` value. If the routine has not been modified since its creation, this column holds the same value as the `CREATED` column.
- `SQL_MODE`: The SQL mode in effect when the routine was created or altered, and under which the routine executes. For the permitted values, see Section 5.1.8, “Server SQL Modes”.
- `CHARACTER_SET_CLIENT`: The session value of the `character_set_client` system variable when the routine was created.
- `COLLATION_CONNECTION`: The session value of the `collation_connection` system variable when the routine was created.
- `DATABASE_COLLATION`: The collation of the database with which the routine is associated.
- The `DATA_TYPE`, `CHARACTER_MAXIMUM_LENGTH`, `CHARACTER_OCTET_LENGTH`, `NUMERIC_PRECISION`, `NUMERIC_SCALE`, `DATETIME_PRECISION`, `CHARACTER_SET_NAME`, and `COLLATION_NAME` columns provide information about the data type for the `RETURNS` clause of stored functions. If a stored routine is a stored procedure, these columns all are `NULL`.
- Information about stored function `RETURNS` data types is also available in the `PARAMETERS` table. The return value data type row for a function can be identified as the row that has an `ORDINAL_POSITION` value of 0.

24.19 The INFORMATION_SCHEMA SCHEMATA Table

A schema is a database, so the `SCHEMATA` table provides information about databases.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`CATALOG_NAME`		`def`
`SCHEMA_NAME`	Database	
`DEFAULT_CHARACTER_SET_NAME`		
`DEFAULT_COLLATION_NAME`		
`SQL_PATH`		`NULL`

The following statements are equivalent:

```
SELECT SCHEMA_NAME AS `Database`
  FROM INFORMATION_SCHEMA.SCHEMATA
  [WHERE SCHEMA_NAME LIKE 'wild']

SHOW DATABASES
  [LIKE 'wild']
```

24.20 The INFORMATION_SCHEMA SCHEMA_PRIVILEGES Table

The `SCHEMA_PRIVILEGES` table provides information about schema (database) privileges. This information comes from the `mysql.db` grant table.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`GRANTEE`		`'user_name'@'host_name'` value, MySQL extension
`TABLE_CATALOG`		`def`, MySQL extension
`TABLE_SCHEMA`		MySQL extension

INFORMATION_SCHEMA Name	SHOW Name	Remarks
PRIVILEGE_TYPE		MySQL extension
IS_GRANTABLE		MySQL extension

Notes:

- This is a nonstandard table. It takes its values from the `mysql.db` table.

24.21 The INFORMATION_SCHEMA STATISTICS Table

The `STATISTICS` table provides information about table indexes.

Columns in `STATISTICS` that represent table statistics hold cached values. If the `information_schema_stats` system variable is `CACHED` (the default), the server retrieves the cached values for queries on the `STATISTICS` table. (To update the cached values for a given table, use `ANALYZE TABLE`.) To use the latest statistics directly from storage engines, set `information_schema_stats` to `LATEST`. For more information, see Section 8.2.3, "Optimizing INFORMATION_SCHEMA Queries".

Note

If the `innodb_read_only` system variable is enabled, `ANALYZE TABLE` may fail because it cannot update statistics tables in the data dictionary, which use `InnoDB`. For `ANALYZE TABLE` operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a `MyISAM` table). To obtain the updated distribution statistics, set `information_schema_stats=LATEST`.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLE_CATALOG		def
TABLE_SCHEMA		= Database
TABLE_NAME	Table	
NON_UNIQUE	Non_unique	
INDEX_SCHEMA		= Database
INDEX_NAME	Key_name	
SEQ_IN_INDEX	Seq_in_index	
COLUMN_NAME	Column_name	
COLLATION	Collation	
CARDINALITY	Cardinality	
SUB_PART	Sub_part	MySQL extension
PACKED	Packed	MySQL extension
NULLABLE	Null	MySQL extension
INDEX_TYPE	Index_type	MySQL extension
COMMENT	Comment	MySQL extension
INDEX_COMMENT	Index_comment	MySQL extension
IS_VISIBLE	Visible	MySQL extension

Notes:

- There is no standard table for indexes. The preceding list is similar to what SQL Server 2000 returns for `sp_statistics`, except that we replaced the name `QUALIFIER` with `CATALOG` and we replaced the name `OWNER` with `SCHEMA`.

Clearly, the preceding table and the output from `SHOW INDEX` are derived from the same parent. So the correlation is already close.

The following statements are equivalent:

```
SELECT * FROM INFORMATION_SCHEMA.STATISTICS
  WHERE table_name = 'tbl_name'
  AND table_schema = 'db_name'

SHOW INDEX
  FROM tbl_name
  FROM db_name
```

24.22 The INFORMATION_SCHEMA ST_GEOMETRY_COLUMNS Table

The `ST_GEOMETRY_COLUMNS` table provides information about table columns that store spatial data. This table is based on the SQL/MM (ISO/IEC 13249-3) standard, with extensions as noted. MySQL implements `ST_GEOMETRY_COLUMNS` as a view on the `INFORMATION_SCHEMA COLUMNS` table.

`ST_GEOMETRY_COLUMNS` has these columns:

- `TABLE_CATALOG`

  The `COLUMNS.TABLE_CATALOG` value.

- `TABLE_SCHEMA`

  The `COLUMNS.TABLE_SCHEMA` value.

- `TABLE_NAME`

  The `COLUMNS.TABLE_NAME` value.

- `COLUMN_NAME`

  The `COLUMNS.TABLE_COLUMN_NAME` value.

- `SRS_NAME`

  Reserved for future use. Currently always `NULL`.

- `SRS_ID`

  Reserved for future use. Currently always `NULL`.

- `GEOMETRY_TYPE_NAME`

  The column data type. Permitted values are: `geometry`, `point`, `linestring`, `polygon`, `multipoint`, `multilinestring`, `multipolygon`, `geometrycollection`. This column is a MySQL extension to the standard.

24.23 The INFORMATION_SCHEMA ST_SPATIAL_REFERENCE_SYSTEMS Table

The `ST_SPATIAL_REFERENCE_SYSTEMS` table provides information about available spatial reference systems for spatial data. This table is based on the SQL/MM (ISO/IEC 13249-3) standard. MySQL implements `ST_SPATIAL_REFERENCE_SYSTEMS` as a view on the `st_spatial_reference_systems` data dictionary table.

ST_SPATIAL_REFERENCE_SYSTEMS has these columns:

- SRS_NAME

  The spatial reference system name. This value is unique.

- SRS_ID

  The spatial reference system numeric ID. This value is unique.

  SRS_ID values represent the same kind of values passed as the SRID argument to spatial functions.

  SRID 0 is special. It is always a legal spatial reference system ID and can be used in any computations on spatial data that depend on SRID values. However, SRID 0 has no entry in ST_SPATIAL_REFERENCE_SYSTEMS because it cannot be described by the SRS definition format. It represents an infinite flat Cartesian plane with no units assigned to its axes. Unlike projected SRSs, it is not georeferenced and it does not necessarily represent Earth. It is an abstract plane that can be used for anything.

- ORGANIZATION

  The name of the organization that defined the coordinate system on which the spatial reference system is based.

- ORGANIZATION_COORDSYS_ID

  The numeric ID given to the spatial reference system by the organization that defined it.

- DEFINITION

  The spatial reference system definition, as a WKT value.

- DESCRIPTION

  The spatial reference system description.

Notes:

- The SRS_NAME, ORGANIZATION, ORGANIZATION_COORDSYS_ID, and DESCRIPTION columns contain information that may be of interest to users, but they are not used by MySQL.

Example:

```
mysql> SELECT * FROM ST_SPATIAL_REFERENCE_SYSTEMS
       WHERE SRS_ID = 4326\G
*************************** 1. row ***************************
               SRS_NAME: WGS 84
                 SRS_ID: 4326
           ORGANIZATION: EPSG
ORGANIZATION_COORDSYS_ID: 4326
             DEFINITION: GEOGCS["WGS 84",DATUM["World Geodetic System 1984",
                         SPHEROID["WGS 84",6378137,298.257223563,
                         AUTHORITY["EPSG","7030"]],AUTHORITY["EPSG","6326"]],
                         PRIMEM["Greenwich",0,AUTHORITY["EPSG","8901"]],
                         UNIT["degree",0.017453292519943278,
                         AUTHORITY["EPSG","9122"]],
                         AXIS["Lat",NORTH],AXIS["Long",EAST],
                         AUTHORITY["EPSG","4326"]]
            DESCRIPTION:
```

24.24 The INFORMATION_SCHEMA TABLES Table

The TABLES table provides information about tables in databases.

Columns in TABLES that represent table statistics hold cached values. If the information_schema_stats system variable is CACHED (the default), the server retrieves the cached values for queries on the TABLES table. (To update the cached values for a given table, use ANALYZE TABLE.) To use the latest statistics directly from storage engines, set information_schema_stats to LATEST. For more information, see Section 8.2.3, "Optimizing INFORMATION_SCHEMA Queries".

Note

If the innodb_read_only system variable is enabled, ANALYZE TABLE may fail because it cannot update statistics tables in the data dictionary, which use InnoDB. For ANALYZE TABLE operations that update the key distribution, failure may occur even if the operation updates the table itself (for example, if it is a MyISAM table). To obtain the updated distribution statistics, set information_schema_stats=LATEST.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLE_CATALOG		def
TABLE_SCHEMA	Table_...	
TABLE_NAME	Table_...	
TABLE_TYPE		
ENGINE	Engine	MySQL extension
VERSION	Version	The .frm file version. With the removal of .frm files in MySQL 8.0, this field now reports a hardcoded value of 10, which is the last .frm file version used in MySQL 5.7.
ROW_FORMAT	Row_format	MySQL extension
TABLE_ROWS	Rows	MySQL extension
AVG_ROW_LENGTH	Avg_row_length	MySQL extension
DATA_LENGTH	Data_length	MySQL extension
MAX_DATA_LENGTH	Max_data_length	MySQL extension
INDEX_LENGTH	Index_length	MySQL extension
DATA_FREE	Data_free	MySQL extension
AUTO_INCREMENT	Auto_increment	MySQL extension
CREATE_TIME	Create_time	MySQL extension
UPDATE_TIME	Update_time	MySQL extension
CHECK_TIME	Check_time	MySQL extension
TABLE_COLLATION	Collation	MySQL extension
CHECKSUM	Checksum	MySQL extension
CREATE_OPTIONS	Create_options	MySQL extension
TABLE_COMMENT	Comment	MySQL extension

Notes:

- Refer to SHOW TABLE STATUS for field descriptions.
- TABLE_SCHEMA and TABLE_NAME are a single field in a SHOW display, for example Table_in_db1.
- TABLE_TYPE should be BASE TABLE or VIEW. The TABLES table does not list TEMPORARY tables.

- For partitioned tables, the ENGINE column shows the name of the storage engine used by all partitions. (Previously, this column showed PARTITION for such tables.)
- The TABLE_ROWS column is NULL if the table is in the INFORMATION_SCHEMA database.

  For InnoDB tables, the row count is only a rough estimate used in SQL optimization. (This is also true if the InnoDB table is partitioned.)
- The DATA_FREE column shows the free space in bytes for InnoDB tables.
- UPDATE_TIME displays a timestamp value for the last UPDATE, INSERT, or DELETE performed on InnoDB tables that are not partitioned. For MVCC, the timestamp value reflects the COMMIT time, which is considered the last update time. Timestamps are not persisted when the server is restarted or when the table is evicted from the InnoDB data dictionary cache.
- For partitioned InnoDB tables, the CHECK_TIME column is always NULL.
- We have nothing for the table's default character set. TABLE_COLLATION is close, because collation names begin with a character set name.
- The CREATE_OPTIONS column shows partitioned if the table is partitioned.

The following statements are equivalent:

```
SELECT table_name FROM INFORMATION_SCHEMA.TABLES
  WHERE table_schema = 'db_name'
  [AND table_name LIKE 'wild']

SHOW TABLES
  FROM db_name
  [LIKE 'wild']
```

24.25 The INFORMATION_SCHEMA TABLESPACES Table

The TABLESPACES table provides information about active tablespaces.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLESPACE_NAME		MySQL extension
ENGINE		MySQL extension
TABLESPACE_TYPE		MySQL extension
LOGFILE_GROUP_NAME		MySQL extension
EXTENT_SIZE		MySQL extension
AUTOEXTEND_SIZE		MySQL extension
MAXIMUM_SIZE		MySQL extension
NODEGROUP_ID		MySQL extension
TABLESPACE_COMMENT		MySQL extension

Notes:

The INFORMATION_SCHEMA.TABLESPACES table does not provide information about InnoDB tablespaces. For InnoDB tablespace metadata, see INNODB_SYS_TABLESPACES and INNODB_SYS_DATAFILES. The INFORMATION_SCHEMA.FILES table also provides metadata for InnoDB tablespaces.

24.26 The INFORMATION_SCHEMA TABLE_CONSTRAINTS Table

The TABLE_CONSTRAINTS table describes which tables have constraints.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
CONSTRAINT_CATALOG		def
CONSTRAINT_SCHEMA		
CONSTRAINT_NAME		
TABLE_SCHEMA		
TABLE_NAME		
CONSTRAINT_TYPE		

Notes:

- The CONSTRAINT_TYPE value can be UNIQUE, PRIMARY KEY, or FOREIGN KEY.
- The UNIQUE and PRIMARY KEY information is about the same as what you get from the Key_name field in the output from SHOW INDEX when the Non_unique field is 0.
- The CONSTRAINT_TYPE column can contain one of these values: UNIQUE, PRIMARY KEY, FOREIGN KEY, CHECK. This is a CHAR (not ENUM) column. The CHECK value is not available until we support CHECK.

24.27 The INFORMATION_SCHEMA TABLE_PRIVILEGES Table

The TABLE_PRIVILEGES table provides information about table privileges. This information comes from the mysql.tables_priv grant table.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
GRANTEE		*'user_name'@'host_name'* value
TABLE_CATALOG		def
TABLE_SCHEMA		
TABLE_NAME		
PRIVILEGE_TYPE		
IS_GRANTABLE		

Notes:

- PRIVILEGE_TYPE can contain one (and only one) of these values: SELECT, INSERT, UPDATE, REFERENCES, ALTER, INDEX, DROP, CREATE VIEW.

The following statements are *not* equivalent:

```
SELECT ... FROM INFORMATION_SCHEMA.TABLE_PRIVILEGES

SHOW GRANTS ...
```

24.28 The INFORMATION_SCHEMA TRIGGERS Table

The TRIGGERS table provides information about triggers. To see information about a table's triggers, you must have the TRIGGER privilege for the table.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TRIGGER_CATALOG		def
TRIGGER_SCHEMA		
TRIGGER_NAME	Trigger	
EVENT_MANIPULATION	Event	

INFORMATION_SCHEMA Name	SHOW Name	Remarks
EVENT_OBJECT_CATALOG		def
EVENT_OBJECT_SCHEMA		
EVENT_OBJECT_TABLE	Table	
ACTION_ORDER		
ACTION_CONDITION		NULL
ACTION_STATEMENT	Statement	
ACTION_ORIENTATION		ROW
ACTION_TIMING	Timing	
ACTION_REFERENCE_OLD_TABLE		NULL
ACTION_REFERENCE_NEW_TABLE		NULL
ACTION_REFERENCE_OLD_ROW		OLD
ACTION_REFERENCE_NEW_ROW		NEW
CREATED	Created	
SQL_MODE	sql_mode	MySQL extension
DEFINER	Definer	MySQL extension
CHARACTER_SET_CLIENT	character_set_client	MySQL extension
COLLATION_CONNECTION	collation_connection	MySQL extension
DATABASE_COLLATION	Database Collation	MySQL extension

Notes:

- The names in the "SHOW Name" column refer to the SHOW TRIGGERS statement, not SHOW CREATE TRIGGER. See Section 13.7.5.38, "SHOW TRIGGERS Syntax".
- TRIGGER_SCHEMA and TRIGGER_NAME: The name of the database in which the trigger occurs and the trigger name, respectively.
- EVENT_MANIPULATION: The trigger event. This is the type of operation on the associated table for which the trigger activates. The value is 'INSERT' (a row was inserted), 'DELETE' (a row was deleted), or 'UPDATE' (a row was modified).
- EVENT_OBJECT_SCHEMA and EVENT_OBJECT_TABLE: As noted in Section 23.3, "Using Triggers", every trigger is associated with exactly one table. These columns indicate the database in which this table occurs, and the table name, respectively.
- ACTION_ORDER: The ordinal position of the trigger's action within the list of triggers on the same table with the same EVENT_MANIPULATION and ACTION_TIMING values.
- ACTION_STATEMENT: The trigger body; that is, the statement executed when the trigger activates. This text uses UTF-8 encoding.
- ACTION_ORIENTATION: Always contains the value 'ROW'.
- ACTION_TIMING: Whether the trigger activates before or after the triggering event. The value is 'BEFORE' or 'AFTER'.
- ACTION_REFERENCE_OLD_ROW and ACTION_REFERENCE_NEW_ROW: The old and new column identifiers, respectively. This means that ACTION_REFERENCE_OLD_ROW always contains the value 'OLD' and ACTION_REFERENCE_NEW_ROW always contains the value 'NEW'.
- CREATED: The date and time when the trigger was created. This is a TIMESTAMP(2) value (with a fractional part in hundredths of seconds) for triggers.

- `SQL_MODE`: The SQL mode in effect when the trigger was created, and under which the trigger executes. For the permitted values, see Section 5.1.8, “Server SQL Modes”.
- `DEFINER`: The account of the user who created the trigger, in `'user_name'@'host_name'` format.
- `CHARACTER_SET_CLIENT`: The session value of the `character_set_client` system variable when the trigger was created.
- `COLLATION_CONNECTION`: The session value of the `collation_connection` system variable when the trigger was created.
- `DATABASE_COLLATION`: The collation of the database with which the trigger is associated.
- The following columns currently always contain `NULL`: `ACTION_CONDITION`, `ACTION_REFERENCE_OLD_TABLE`, and `ACTION_REFERENCE_NEW_TABLE`.

Example, using the `ins_sum` trigger defined in Section 23.3, “Using Triggers”:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.TRIGGERS
       WHERE TRIGGER_SCHEMA='test' AND TRIGGER_NAME='ins_sum'\G
*************************** 1. row ***************************
           TRIGGER_CATALOG: def
            TRIGGER_SCHEMA: test
              TRIGGER_NAME: ins_sum
        EVENT_MANIPULATION: INSERT
      EVENT_OBJECT_CATALOG: def
       EVENT_OBJECT_SCHEMA: test
        EVENT_OBJECT_TABLE: account
              ACTION_ORDER: 1
          ACTION_CONDITION: NULL
          ACTION_STATEMENT: SET @sum = @sum + NEW.amount
        ACTION_ORIENTATION: ROW
             ACTION_TIMING: BEFORE
ACTION_REFERENCE_OLD_TABLE: NULL
ACTION_REFERENCE_NEW_TABLE: NULL
  ACTION_REFERENCE_OLD_ROW: OLD
  ACTION_REFERENCE_NEW_ROW: NEW
                   CREATED: 2017-03-30 20:29:39.06
                 SQL_MODE: ONLY_FULL_GROUP_BY,STRICT_TRANS_TABLES,NO_ZERO_IN_DATE,NO_ZERO_DATE,ERROR_F
                  DEFINER: me@localhost
      CHARACTER_SET_CLIENT: utf8
      COLLATION_CONNECTION: utf8_general_ci
        DATABASE_COLLATION: utf8mb4_0900_ai_ci
```

24.29 The INFORMATION_SCHEMA USER_PRIVILEGES Table

The `USER_PRIVILEGES` table provides information about global privileges. This information comes from the `mysql.user` grant table.

`INFORMATION_SCHEMA` Name	`SHOW` Name	Remarks
`GRANTEE`		`'user_name'@'host_name'` value, MySQL extension
`TABLE_CATALOG`		`def`, MySQL extension
`PRIVILEGE_TYPE`		MySQL extension
`IS_GRANTABLE`		MySQL extension

Notes:

- This is a nonstandard table. It takes its values from the `mysql.user` table.

24.30 The INFORMATION_SCHEMA VIEWS Table

The VIEWS table provides information about views in databases. You must have the SHOW VIEW privilege to access this table.

INFORMATION_SCHEMA Name	SHOW Name	Remarks
TABLE_CATALOG		def
TABLE_SCHEMA		
TABLE_NAME		
VIEW_DEFINITION		
CHECK_OPTION		
IS_UPDATABLE		
DEFINER		
SECURITY_TYPE		
CHARACTER_SET_CLIENT		MySQL extension
COLLATION_CONNECTION		MySQL extension

Notes:

- The VIEW_DEFINITION column has most of what you see in the Create Table field that SHOW CREATE VIEW produces. Skip the words before SELECT and skip the words WITH CHECK OPTION. Suppose that the original statement was:

  ```
  CREATE VIEW v AS
    SELECT s2,s1 FROM t
    WHERE s1 > 5
    ORDER BY s1
    WITH CHECK OPTION;
  ```

  Then the view definition looks like this:

  ```
  SELECT s2,s1 FROM t WHERE s1 > 5 ORDER BY s1
  ```

- The CHECK_OPTION column has a value of NONE, CASCADE, or LOCAL.

- MySQL sets a flag, called the view updatability flag, at CREATE VIEW time. The flag is set to YES (true) if UPDATE and DELETE (and similar operations) are legal for the view. Otherwise, the flag is set to NO (false). The IS_UPDATABLE column in the VIEWS table displays the status of this flag. It means that the server always knows whether a view is updatable.

  If a view is not updatable, statements such UPDATE, DELETE, and INSERT are illegal and will be rejected. (Note that even if a view is updatable, it might not be possible to insert into it; for details, refer to Section 23.5.3, “Updatable and Insertable Views”.)

- DEFINER: The account of the user who created the view, in *'user_name'@'host_name'* format. SECURITY_TYPE has a value of DEFINER or INVOKER.

- CHARACTER_SET_CLIENT: The session value of the character_set_client system variable when the view was created.

- COLLATION_CONNECTION: The session value of the collation_connection system variable when the view was created.

MySQL lets you use different sql_mode settings to tell the server the type of SQL syntax to support. For example, you might use the ANSI SQL mode to ensure MySQL correctly interprets the standard SQL concatenation operator, the double bar (||), in your queries. If you then create a view that concatenates items, you might worry that changing the sql_mode setting to a value different from ANSI could cause the view to become invalid. But this is not the case. No matter how you write out a

view definition, MySQL always stores it the same way, in a canonical form. Here is an example that shows how the server changes a double bar concatenation operator to a `CONCAT()` function:

```
mysql> SET sql_mode = 'ANSI';
Query OK, 0 rows affected (0.00 sec)

mysql> CREATE VIEW test.v AS SELECT 'a' || 'b' as col1;
Query OK, 0 rows affected (0.00 sec)

mysql> SELECT VIEW_DEFINITION FROM INFORMATION_SCHEMA.VIEWS
       WHERE TABLE_SCHEMA = 'test' AND TABLE_NAME = 'v';
+--------------------------------+
| VIEW_DEFINITION                |
+--------------------------------+
| select concat('a','b') AS `col1` |
+--------------------------------+
1 row in set (0.00 sec)
```

The advantage of storing a view definition in canonical form is that changes made later to the value of `sql_mode` will not affect the results from the view. However an additional consequence is that comments prior to `SELECT` are stripped from the definition by the server.

24.31 InnoDB INFORMATION_SCHEMA Tables

This section provides table definitions for `InnoDB INFORMATION_SCHEMA` tables. For related information and examples, see Section 15.14, “InnoDB INFORMATION_SCHEMA Tables”.

`InnoDB INFORMATION_SCHEMA` tables can be used to monitor ongoing `InnoDB` activity, to detect inefficiencies before they turn into issues, or to troubleshoot performance and capacity issues. As your database becomes bigger and busier, running up against the limits of your hardware capacity, you monitor and tune these aspects to keep the database running smoothly.

24.31.1 The INFORMATION_SCHEMA INNODB_BUFFER_PAGE Table

The `INNODB_BUFFER_PAGE` table holds information about each page in the `InnoDB` buffer pool.

For related usage information and examples, see Section 15.14.5, “InnoDB INFORMATION_SCHEMA Buffer Pool Tables”.

Warning

Querying the `INNODB_BUFFER_PAGE` table can introduce significant performance overhead. Do not query this table on a production system unless you are aware of the performance impact that your query may have, and have determined it to be acceptable. To avoid impacting performance, reproduce the issue you want to investigate on a test instance and query the `INNODB_BUFFER_PAGE` table on the test instance.

Table 24.1 INNODB_BUFFER_PAGE Columns

Column name	Description
`POOL_ID`	Buffer Pool ID. An identifier to distinguish between multiple buffer pool instances.
`BLOCK_ID`	Buffer Pool Block ID.
`SPACE`	Tablespace ID. Uses the same value as in `INNODB_SYS_TABLES.SPACE`.
`PAGE_NUMBER`	Page number.
`PAGE_TYPE`	Page type. Permitted values are `ALLOCATED` (Freshly allocated page), `INDEX` (B-tree node), `UNDO_LOG` (Undo log page), `INODE` (Index

Column name	Description
	node), IBUF_FREE_LIST (Insert buffer free list), IBUF_BITMAP (Insert buffer bitmap), SYSTEM (System page), TRX_SYSTEM (Transaction system data), FILE_SPACE_HEADER (File space header), EXTENT_DESCRIPTOR (Extent descriptor page), BLOB (Uncompressed BLOB page), COMPRESSED_BLOB (First compressed BLOB page), COMPRESSED_BLOB2 (Subsequent comp BLOB page), IBUF_INDEX (Insert buffer index), RTREE_INDEX (spatial index), SDI_INDEX (SDI index), SDI_BLOB (uncompressed SDI BLOB), SDI_COMPRESSED_BLOB (compressed SDI BLOB), UNKNOWN (unknown).
FLUSH_TYPE	Flush type.
FIX_COUNT	Number of threads using this block within the buffer pool. When zero, the block is eligible to be evicted.
IS_HASHED	Whether hash index has been built on this page.
NEWEST_MODIFICATION	Log Sequence Number of the youngest modification.
OLDEST_MODIFICATION	Log Sequence Number of the oldest modification.
ACCESS_TIME	An abstract number used to judge the first access time of the page.
TABLE_NAME	Name of the table the page belongs to. This column is only applicable to pages of type INDEX.
INDEX_NAME	Name of the index the page belongs to. It can be the name of a clustered index or a secondary index. This column is only applicable to pages of type INDEX.
NUMBER_RECORDS	Number of records within the page.
DATA_SIZE	Sum of the sizes of the records. This column is only applicable to pages of type INDEX.
COMPRESSED_SIZE	Compressed page size. Null for pages that are not compressed.
PAGE_STATE	Page state. A page with valid data has one of the following states: FILE_PAGE (buffers a page of data from a file), MEMORY (buffers a page from an in-memory object), COMPRESSED. Other possible states (managed by InnoDB) are: NULL, READY_FOR_USE, NOT_USED, REMOVE_HASH.
IO_FIX	Specifies whether any I/O is pending for this page: IO_NONE = no pending I/O, IO_READ = read pending, IO_WRITE = write pending.
IS_OLD	Specifies whether or not the block is in the sublist of old blocks in the LRU list.
FREE_PAGE_CLOCK	The value of the freed_page_clock counter when the block was the last placed at the head of the LRU list. The freed_page_clock counter tracks the number of blocks removed from the end of the LRU list.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE LIMIT 1\G
*************************** 1. row ***************************
           POOL_ID: 0
         BLOCK_ID: 0
            SPACE: 97
      PAGE_NUMBER: 2473
        PAGE_TYPE: INDEX
       FLUSH_TYPE: 1
        FIX_COUNT: 0
        IS_HASHED: YES
NEWEST_MODIFICATION: 733855581
OLDEST_MODIFICATION: 0
```

```
       ACCESS_TIME: 3378385672
        TABLE_NAME: `employees`.`salaries`
        INDEX_NAME: PRIMARY
    NUMBER_RECORDS: 468
         DATA_SIZE: 14976
   COMPRESSED_SIZE: 0
        PAGE_STATE: FILE_PAGE
            IO_FIX: IO_NONE
            IS_OLD: YES
   FREE_PAGE_CLOCK: 66
```

Notes:

- This table is primarily useful for expert-level performance monitoring, or when developing performance-related extensions for MySQL.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- When tables, table rows, partitions, or indexes are deleted, associated pages remain in the buffer pool until space is required for other data. The `INNODB_BUFFER_PAGE` table reports information about these pages until they are evicted from the buffer pool. For more information about how the `InnoDB` manages buffer pool data, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

24.31.2 The INFORMATION_SCHEMA INNODB_BUFFER_PAGE_LRU Table

The `INNODB_BUFFER_PAGE_LRU` table holds information about the pages in the `InnoDB` buffer pool, in particular how they are ordered in the LRU list that determines which pages to evict from the buffer pool when it becomes full.

The `INNODB_BUFFER_PAGE_LRU` table has the same columns as the `INNODB_BUFFER_PAGE` table, except that the `INNODB_BUFFER_PAGE_LRU` table has an `LRU_POSITION` column instead of a `BLOCK_ID` column.

For related usage information and examples, see Section 15.14.5, “InnoDB INFORMATION_SCHEMA Buffer Pool Tables”.

Warning

Querying the `INNODB_BUFFER_PAGE_LRU` table can introduce significant performance overhead. Do not query this table on a production system unless you are aware of the performance impact that your query may have, and have determined it to be acceptable. To avoid impacting performance, reproduce the issue you want to investigate on a test instance and query the `INNODB_BUFFER_PAGE_LRU` table on the test instance.

Table 24.2 INNODB_BUFFER_PAGE_LRU Columns

Column name	Description
`POOL_ID`	Buffer Pool ID. An identifier to distinguish between multiple buffer pool instances.
`LRU_POSITION`	The position of the page in the LRU list.
`SPACE`	Tablespace ID. Uses the same value as in `INNODB_SYS_TABLES.SPACE`.
`PAGE_NUMBER`	Page number.
`PAGE_TYPE`	Page type. Permitted values are `ALLOCATED` (Freshly allocated page), `INDEX` (B-tree node), `UNDO_LOG` (Undo log page), `INODE` (Index node), `IBUF_FREE_LIST` (Insert buffer free list), `IBUF_BITMAP`

Column name	Description
	(Insert buffer bitmap), SYSTEM (System page), TRX_SYSTEM (Transaction system data), FILE_SPACE_HEADER (File space header), EXTENT_DESCRIPTOR (Extent descriptor page), BLOB (Uncompressed BLOB page), COMPRESSED_BLOB (First compressed BLOB page), COMPRESSED_BLOB2 (Subsequent comp BLOB page), IBUF_INDEX (Insert buffer index), RTREE_INDEX (spatial index), SDI_INDEX (SDI index), SDI_BLOB (uncompressed SDI BLOB), SDI_COMPRESSED_BLOB (compressed SDI BLOB), UNKNOWN (unknown).
FLUSH_TYPE	Flush type.
FIX_COUNT	Number of threads using this block within the buffer pool. When zero, the block is eligible to be evicted.
IS_HASHED	Whether hash index has been built on this page.
NEWEST_MODIFICATION	Log Sequence Number of the youngest modification.
OLDEST_MODIFICATION	Log Sequence Number of the oldest modification.
ACCESS_TIME	An abstract number used to judge the first access time of the page.
TABLE_NAME	Name of the table the page belongs to. This column is only applicable to pages of type INDEX.
INDEX_NAME	Name of the index the page belongs to. It can be the name of a clustered index or a secondary index. This column is only applicable to pages of type INDEX.
NUMBER_RECORDS	Number of records within the page.
DATA_SIZE	Sum of the sizes of the records. This column is only applicable to pages of type INDEX.
COMPRESSED_SIZE	Compressed page size. Null for pages that are not compressed.
PAGE_STATE	Page state. A page with valid data has one of the following states: FILE_PAGE (buffers a page of data from a file), MEMORY (buffers a page from an in-memory object), COMPRESSED. Other possible states (managed by InnoDB) are: NULL, READY_FOR_USE, NOT_USED, REMOVE_HASH.
IO_FIX	Specifies whether any I/O is pending for this page: IO_NONE = no pending I/O, IO_READ = read pending, IO_WRITE = write pending.
IS_OLD	Specifies whether or not the block is in the sublist of old blocks in the LRU list.
FREE_PAGE_CLOCK	The value of the freed_page_clock counter when the block was the last placed at the head of the LRU list. The freed_page_clock counter tracks the number of blocks removed from the end of the LRU list.

Example

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_BUFFER_PAGE_LRU LIMIT 1\G
*************************** 1. row ***************************
            POOL_ID: 0
      LRU_POSITION: 0
             SPACE: 97
       PAGE_NUMBER: 1984
         PAGE_TYPE: INDEX
        FLUSH_TYPE: 1
         FIX_COUNT: 0
         IS_HASHED: YES
NEWEST_MODIFICATION: 719490396
OLDEST_MODIFICATION: 0
       ACCESS_TIME: 3378383796
```

```
        TABLE_NAME: `employees`.`salaries`
        INDEX_NAME: PRIMARY
    NUMBER_RECORDS: 468
         DATA_SIZE: 14976
   COMPRESSED_SIZE: 0
        COMPRESSED: NO
            IO_FIX: IO_NONE
            IS_OLD: YES
   FREE_PAGE_CLOCK: 0
```

Notes

- This table is primarily useful for expert-level performance monitoring, or when developing performance-related extensions for MySQL.
- You must have the `PROCESS` privilege to query this table.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- Querying this table can require MySQL to allocate a large block of contiguous memory, more than 64 bytes time the number of active pages in the buffer pool. This allocation could potentially cause an out-of-memory error, especially for systems with multi-gigabyte buffer pools.
- Querying this table requires MySQL to lock the data structure representing the buffer pool while traversing the LRU list, which can reduce concurrency, especially for systems with multi-gigabyte buffer pools.
- When tables, table rows, partitions, or indexes are deleted, associated pages remain in the buffer pool until space is required for other data. The `INNODB_BUFFER_PAGE_LRU` table reports information about these pages until they are evicted from the buffer pool. For more information about how the `InnoDB` manages buffer pool data, see Section 15.6.3.1, “The InnoDB Buffer Pool”.

24.31.3 The INFORMATION_SCHEMA INNODB_BUFFER_POOL_STATS Table

The `INNODB_BUFFER_POOL_STATS` table provides much of the same buffer pool information provided in `SHOW ENGINE INNODB STATUS` output. Much of the same information may also be obtained using `InnoDB` buffer pool server status variables.

The idea of making pages in the buffer pool “young” or “not young” refers to transferring them between the sublists at the head and tail of the buffer pool data structure. Pages made “young” take longer to age out of the buffer pool, while pages made “not young” are moved much closer to the point of eviction.

For related usage information and examples, see Section 15.14.5, “InnoDB INFORMATION_SCHEMA Buffer Pool Tables”.

Table 24.3 INNODB_BUFFER_POOL_STATS Columns

Column name	Description
`POOL_ID`	Buffer Pool ID. A unique identifier to distinguish between multiple buffer pool instances.
`POOL_SIZE`	The `InnoDB` buffer pool size in pages.
`FREE_BUFFERS`	The number of free pages in the `InnoDB` buffer pool
`DATABASE_PAGES`	The number of pages in the `InnoDB` buffer pool containing data. The number includes both dirty and clean pages.
`OLD_DATABASE_PAGES`	The number of pages in the `old` buffer pool sublist.
`MODIFIED_DATABASE_PAGES`	The number of modified (dirty) database pages

Column name	Description
PENDING_DECOMPRESS	The number of pages pending decompression
PENDING_READS	The number of pending reads
PENDING_FLUSH_LRU	The number of pages pending flush in the LRU
PENDING_FLUSH_LIST	The number of pages pending flush in the flush list
PAGES_MADE_YOUNG	The number of pages made young
PAGES_NOT_MADE_YOUNG	The number of pages not made young
PAGES_MADE_YOUNG_RATE	The number of pages made young per second (pages made young since the last printout / time elapsed)
PAGES_MADE_NOT_YOUNG_RATE	The number of pages not made per second (pages not made young since the last printout / time elapsed)
NUMBER_PAGES_READ	The number of pages read
NUMBER_PAGES_CREATED	The number of pages created
NUMBER_PAGES_WRITTEN	The number of pages written
PAGES_READ_RATE	The number of pages read per second (pages read since the last printout / time elapsed)
PAGES_CREATE_RATE	The number of pages created per second (pages created since the last printout / time elapsed)
PAGES_WRITTEN_RATE	The number of pages written per second (pages written since the last printout / time elapsed)
NUMBER_PAGES_GET	The number of logical read requests.
HIT_RATE	The buffer pool hit rate
YOUNG_MAKE_PER_THOUSAND_GETS	The number of pages made young per thousand gets
NOT_YOUNG_MAKE_PER_THOUSAND_GETS	The number of pages not made young per thousand gets
NUMBER_PAGES_READ_AHEAD	The number of pages read ahead
NUMBER_READ_AHEAD_EVICTED	The number of pages read into the InnoDB buffer pool by the read-ahead background thread that were subsequently evicted without having been accessed by queries.
READ_AHEAD_RATE	The read ahead rate per second (pages read ahead since the last printout / time elapsed)
READ_AHEAD_EVICTED_RATE	The number of read ahead pages evicted without access per second (read ahead pages not accessed since the last printout / time elapsed)
LRU_IO_TOTAL	LRU IO total
LRU_IO_CURRENT	LRU IO for the current interval
UNCOMPRESS_TOTAL	Total number of pages decompressed
UNCOMPRESS_CURRENT	The number of pages decompressed in the current interval

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_BUFFER_POOL_STATS\G
*************************** 1. row ***************************
                     POOL_ID: 0
                   POOL_SIZE: 8192
                FREE_BUFFERS: 1
              DATABASE_PAGES: 8085
          OLD_DATABASE_PAGES: 2964
      MODIFIED_DATABASE_PAGES: 0
           PENDING_DECOMPRESS: 0
                PENDING_READS: 0
            PENDING_FLUSH_LRU: 0
```

```
             PENDING_FLUSH_LIST: 0
               PAGES_MADE_YOUNG: 22821
           PAGES_NOT_MADE_YOUNG: 3544303
          PAGES_MADE_YOUNG_RATE: 357.62602199870594
      PAGES_MADE_NOT_YOUNG_RATE: 0
             NUMBER_PAGES_READ: 2389
          NUMBER_PAGES_CREATED: 12385
          NUMBER_PAGES_WRITTEN: 13111
               PAGES_READ_RATE: 0
             PAGES_CREATE_RATE: 0
            PAGES_WRITTEN_RATE: 0
              NUMBER_PAGES_GET: 33322210
                      HIT_RATE: 1000
   YOUNG_MAKE_PER_THOUSAND_GETS: 18
NOT_YOUNG_MAKE_PER_THOUSAND_GETS: 0
        NUMBER_PAGES_READ_AHEAD: 2024
      NUMBER_READ_AHEAD_EVICTED: 0
                READ_AHEAD_RATE: 0
        READ_AHEAD_EVICTED_RATE: 0
                    LRU_IO_TOTAL: 0
                  LRU_IO_CURRENT: 0
                UNCOMPRESS_TOTAL: 0
              UNCOMPRESS_CURRENT: 0
```

Notes:

- This table is primarily useful for expert-level performance monitoring, or when developing performance-related extensions for MySQL.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.

24.31.4 The INFORMATION_SCHEMA INNODB_CACHED_INDEXES Table

The `INNODB_CACHED_INDEXES` table reports the number of index pages cached in the `InnoDB` buffer pool for each index.

Table 24.4 INNODB_CACHED_INDEXES Columns

Column name	Description
`SPACE_ID`	Tablespace Space ID.
`INDEX_ID`	An identifier for each index that is unique across all the databases in an instance.
`N_CACHED_PAGES`	The number of index pages cached in the `InnoDB` buffer pool.

Examples:

This query returns the number of index pages cached in the `InnoDB` buffer pool for a specific index:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_CACHED_INDEXES WHERE INDEX_ID=65\G
*************************** 1. row ***************************
      INDEX_ID: 65
N_CACHED_PAGES: 45
1 row in set (0.00 sec)
```

This query returns the number of index pages cached in the `InnoDB` buffer pool for each index, and uses the `INNODB_SYS_INDEXES` and `INNODB_SYS_TABLES` to resolve the table name and index name for each `INDEX_ID` value.

```
SELECT
  tables.name AS table_name,
```

```
  indexes.name AS index_name,
  cached.n_cached_pages AS n_cached_pages
FROM
  INFORMATION_SCHEMA.INNODB_CACHED_INDEXES AS cached,
  INFORMATION_SCHEMA.INNODB_SYS_INDEXES AS indexes,
  INFORMATION_SCHEMA.INNODB_SYS_TABLES AS tables
WHERE
  cached.index_id = indexes.index_id
  AND indexes.table_id = tables.table_id;
```

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.

24.31.5 The INFORMATION_SCHEMA INNODB_CMP and INNODB_CMP_RESET Tables

The `INNODB_CMP` and `INNODB_CMP_RESET` tables contain status information on operations related to compressed `InnoDB` tables.

Table 24.5 Columns of INNODB_CMP and INNODB_CMP_RESET

Column name	Description
`PAGE_SIZE`	Compressed page size in bytes.
`COMPRESS_OPS`	Number of times a B-tree page of the size `PAGE_SIZE` has been compressed. Pages are compressed whenever an empty page is created or the space for the uncompressed modification log runs out.
`COMPRESS_OPS_OK`	Number of times a B-tree page of the size `PAGE_SIZE` has been successfully compressed. This count should never exceed `COMPRESS_OPS`.
`COMPRESS_TIME`	Total time in seconds spent in attempts to compress B-tree pages of the size `PAGE_SIZE`.
`UNCOMPRESS_OPS`	Number of times a B-tree page of the size `PAGE_SIZE` has been uncompressed. B-tree pages are uncompressed whenever compression fails or at first access when the uncompressed page does not exist in the buffer pool.
`UNCOMPRESS_TIME`	Total time in seconds spent in uncompressing B-tree pages of the size `PAGE_SIZE`.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_CMP\G
*************************** 1. row ***************************
      page_size: 1024
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 0
uncompress_time: 0
*************************** 2. row ***************************
      page_size: 2048
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 0
uncompress_time: 0
*************************** 3. row ***************************
      page_size: 4096
```

```
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 0
uncompress_time: 0
*************************** 4. row ***************************
      page_size: 8192
   compress_ops: 86955
compress_ops_ok: 81182
  compress_time: 27
 uncompress_ops: 26828
uncompress_time: 5
*************************** 5. row ***************************
      page_size: 16384
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 0
uncompress_time: 0
```

Notes:

- Use these tables to measure the effectiveness of `InnoDB` table compression in your database.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of these tables including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For usage information, see Section 15.9.1.4, “Monitoring InnoDB Table Compression at Runtime” and Section 15.14.1.3, “Using the Compression Information Schema Tables”. For general information about `InnoDB` table compression, see Section 15.9, “InnoDB Table and Page Compression”.

24.31.6 The INFORMATION_SCHEMA INNODB_CMPMEM and INNODB_CMPMEM_RESET Tables

The `INNODB_CMPMEM` and `INNODB_CMPMEM_RESET` tables contain status information on compressed pages within the `InnoDB` buffer pool.

Table 24.6 Columns of INNODB_CMPMEM and INNODB_CMPMEM_RESET

Column name	Description
`PAGE_SIZE`	Block size in bytes. Each record of this table describes blocks of this size.
`BUFFER_POOL_INSTANCE`	A unique identifier for the buffer pool instance.
`PAGES_USED`	Number of blocks of the size `PAGE_SIZE` that are currently in use.
`PAGES_FREE`	Number of blocks of the size `PAGE_SIZE` that are currently available for allocation. This column shows the external fragmentation in the memory pool. Ideally, these numbers should be at most 1.
`RELOCATION_OPS`	Number of times a block of the size `PAGE_SIZE` has been relocated. The buddy system can relocate the allocated “buddy neighbor” of a freed block when it tries to form a bigger freed block. Reading from the table `INNODB_CMPMEM_RESET` resets this count.
`RELOCATION_TIME`	Total time in microseconds spent in relocating blocks of the size `PAGE_SIZE`. Reading from the table `INNODB_CMPMEM_RESET` resets this count.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_CMPMEM\G
```

```
*************************** 1. row ***************************
           page_size: 1024
buffer_pool_instance: 0
          pages_used: 0
          pages_free: 0
      relocation_ops: 0
     relocation_time: 0
*************************** 2. row ***************************
           page_size: 2048
buffer_pool_instance: 0
          pages_used: 0
          pages_free: 0
      relocation_ops: 0
     relocation_time: 0
*************************** 3. row ***************************
           page_size: 4096
buffer_pool_instance: 0
          pages_used: 0
          pages_free: 0
      relocation_ops: 0
     relocation_time: 0
*************************** 4. row ***************************
           page_size: 8192
buffer_pool_instance: 0
          pages_used: 7673
          pages_free: 15
      relocation_ops: 4638
     relocation_time: 0
*************************** 5. row ***************************
           page_size: 16384
buffer_pool_instance: 0
          pages_used: 0
          pages_free: 0
      relocation_ops: 0
     relocation_time: 0
```

Notes:

- Use these tables to measure the effectiveness of `InnoDB` table compression in your database.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of these tables including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For usage information, see Section 15.9.1.4, “Monitoring InnoDB Table Compression at Runtime” and Section 15.14.1.3, “Using the Compression Information Schema Tables”. For general information about `InnoDB` table compression, see Section 15.9, “InnoDB Table and Page Compression”.

24.31.7 The INFORMATION_SCHEMA INNODB_CMP_PER_INDEX and INNODB_CMP_PER_INDEX_RESET Tables

The `INNODB_CMP_PER_INDEX` and `INNODB_CMP_PER_INDEX_RESET` tables contain status information on operations related to compressed `InnoDB` tables and indexes, with separate statistics for each combination of database, table, and index, to help you evaluate the performance and usefulness of compression for specific tables.

For a compressed `InnoDB` table, both the table data and all the secondary indexes are compressed. In this context, the table data is treated as just another index, one that happens to contain all the columns: the clustered index.

Table 24.7 Columns of INNODB_CMP_PER_INDEX and INNODB_CMP_PER_INDEX_RESET

Column name	Description
`DATABASE_NAME`	Database containing the applicable table.

Column name	Description
TABLE_NAME	Table to monitor for compression statistics.
INDEX_NAME	Index to monitor for compression statistics.
COMPRESS_OPS	Number of compression operations attempted. Pages are compressed whenever an empty page is created or the space for the uncompressed modification log runs out.
COMPRESS_OPS_OK	Number of successful compression operations. Subtract from the COMPRESS_OPS value to get the number of compression failures. Divide by the COMPRESS_OPS value to get the percentage of compression failures.
COMPRESS_TIME	Total amount of CPU time, in seconds, used for compressing data in this index.
UNCOMPRESS_OPS	Number of uncompression operations performed. Compressed InnoDB pages are uncompressed whenever compression fails, or the first time a compressed page is accessed in the buffer pool and the uncompressed page does not exist.
UNCOMPRESS_TIME	Total amount of CPU time, in seconds, used for uncompressing data in this index.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_CMP_PER_INDEX\G
*************************** 1. row ***************************
  database_name: employees
     table_name: salaries
     index_name: PRIMARY
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 23451
uncompress_time: 4
*************************** 2. row ***************************
  database_name: employees
     table_name: salaries
     index_name: emp_no
   compress_ops: 0
compress_ops_ok: 0
  compress_time: 0
 uncompress_ops: 1597
uncompress_time: 0
```

Notes:

- Use these tables to measure the effectiveness of InnoDB table compression for specific tables, indexes, or both.
- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of these tables including data types and default values.
- You must have the PROCESS privilege to query these tables.
- Because collecting separate measurements for every index imposes substantial performance overhead, INNODB_CMP_PER_INDEX and INNODB_CMP_PER_INDEX_RESET statistics are not gathered by default. You must enable the innodb_cmp_per_index_enabled configuration option before performing the operations on compressed tables that you want to monitor.
- For usage information, see Section 15.9.1.4, “Monitoring InnoDB Table Compression at Runtime” and Section 15.14.1.3, “Using the Compression Information Schema Tables”. For general information about InnoDB table compression, see Section 15.9, “InnoDB Table and Page Compression”.

24.31.8 The INFORMATION_SCHEMA INNODB_FT_BEING_DELETED Table

The `INNODB_FT_BEING_DELETED` table is a snapshot of the `INNODB_FT_DELETED` table that is only used during an `OPTIMIZE TABLE` maintenance operation. When `OPTIMIZE TABLE` is run, the `INNODB_FT_BEING_DELETED` table is emptied, and DOC_IDs are removed from the `INNODB_FT_DELETED` table. Because the contents of `INNODB_FT_BEING_DELETED` typically have a short lifetime, this table has limited utility for monitoring or debugging. For information about running `OPTIMIZE TABLE` on tables with `FULLTEXT` indexes, see Section 12.9.6, "Fine-Tuning MySQL Full-Text Search".

This table initially appears empty, until you set the value of the configuration variable `innodb_ft_aux_table`. The output appears similar to the example provided for the `INNODB_FT_DELETED` table.

For related usage information and examples, see Section 15.14.4, "InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables".

Table 24.8 INNODB_FT_BEING_DELETED Columns

Column name	Description
`DOC_ID`	The document ID of the row that is in the process of being deleted. This value might reflect the value of an ID column that you defined for the underlying table, or it can be a sequence value generated by `InnoDB` when the table does not contain a suitable column. This value is used to skip rows in the `innodb_ft_index_table` table, when you do text searches before data for deleted rows is physically removed from the `FULLTEXT` index by an `OPTIMIZE TABLE` statement. See Optimizing InnoDB Full-Text Indexes for more information.

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For more information about `InnoDB FULLTEXT` search, see Section 15.8.2.4, "InnoDB FULLTEXT Indexes", and Section 12.9, "Full-Text Search Functions".

24.31.9 The INFORMATION_SCHEMA INNODB_FT_CONFIG Table

The `INNODB_FT_CONFIG` table displays metadata about the `FULLTEXT` index and associated processing for an `InnoDB` table.

Before you query this table, set the configuration variable `innodb_ft_aux_table` to the name (including the database name) of the table that contains the `FULLTEXT` index, for example `test/articles`.

For related usage information and examples, see Section 15.14.4, "InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables".

Table 24.9 INNODB_FT_CONFIG Columns

Column name	Description
`KEY`	The name designating an item of metadata for an `InnoDB` table containing a `FULLTEXT` index.
`VALUE`	The value associated with the corresponding `KEY` column, reflecting some limit or current value for an aspect of a `FULLTEXT` index for an `InnoDB` table.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_CONFIG;
+---------------------------+-------------------+
| KEY                       | VALUE             |
+---------------------------+-------------------+
| optimize_checkpoint_limit | 180               |
| synced_doc_id             | 0                 |
| stopword_table_name       | test/my_stopwords |
| use_stopword              | 1                 |
+---------------------------+-------------------+
```

Notes:

- This table is only intended for internal configuration. It is not intended for statistical information purposes.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- The values for the `KEY` column might evolve depending on the needs for performance tuning and debugging for `InnoDB` full-text processing. The key values include:
  - `optimize_checkpoint_limit`: The number of seconds after which an `OPTIMIZE TABLE` run will stop.
  - `synced_doc_id`: The next `DOC_ID` to be issued.
  - `stopword_table_name`: The *`database/table`* name for a user defined stopword table. This field appears empty if there is no user-defined stopword table.
  - `use_stopword`: Indicates whether or not a stopword table is used, which is defined when the `FULLTEXT` index is created.
- For more information about `InnoDB FULLTEXT` search, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and Section 12.9, “Full-Text Search Functions”.

24.31.10 The INFORMATION_SCHEMA INNODB_FT_DEFAULT_STOPWORD Table

The `INNODB_FT_DEFAULT_STOPWORD` table holds a list of stopwords that are used by default when creating a `FULLTEXT` index on an `InnoDB` table. For information about the default `InnoDB` stopword list and how to define your own stopword lists, see Section 12.9.4, “Full-Text Stopwords”.

For related usage information and examples, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

Table 24.10 INNODB_FT_DEFAULT_STOPWORD Columns

Column name	Description
`value`	A word that is used by default as a stopword for `FULLTEXT` indexes on `InnoDB` tables. Not used if you override the default stopword processing with either the `innodb_ft_server_stopword_table` or the `innodb_ft_user_stopword_table` option.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_DEFAULT_STOPWORD;
```

```
+-------+
| value |
+-------+
| a     |
| about |
| an    |
| are   |
| as    |
| at    |
| be    |
| by    |
| com   |
| de    |
| en    |
| for   |
| from  |
| how   |
| i     |
| in    |
| is    |
| it    |
| la    |
| of    |
| on    |
| or    |
| that  |
| the   |
| this  |
| to    |
| was   |
| what  |
| when  |
| where |
| who   |
| will  |
| with  |
| und   |
| the   |
| www   |
+-------+
36 rows in set (0.00 sec)
```

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For more information about `InnoDB FULLTEXT` search, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and Section 12.9, “Full-Text Search Functions”.

24.31.11 The INFORMATION_SCHEMA INNODB_FT_DELETED Table

The `INNODB_FT_DELETED` table records rows that are deleted from the `FULLTEXT` index for an `InnoDB` table. To avoid expensive index reorganization during DML operations for an `InnoDB FULLTEXT` index, the information about newly deleted words is stored separately, filtered out of search results when you do a text search, and removed from the main search index only when you issue the `OPTIMIZE TABLE` statement for the `InnoDB` table. See Optimizing InnoDB Full-Text Indexes for more information.

This table initially appears empty, until you set the value of the configuration variable `innodb_ft_aux_table` to the name (including the database name) of the table that contains the `FULLTEXT` index, for example `test/articles`.

For related usage information and examples, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

Table 24.11 INNODB_FT_DELETED Columns

Column name	Description
DOC_ID	The document ID of the newly deleted row. This value might reflect the value of an ID column that you defined for the underlying table, or it can be a sequence value generated by InnoDB when the table does not contain a suitable column. This value is used to skip rows in the innodb_ft_index_table table, when you do text searches before data for deleted rows is physically removed from the FULLTEXT index by an OPTIMIZE TABLE statement. See Optimizing InnoDB Full-Text Indexes for more information.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_FT_DELETED;
+--------+
| DOC_ID |
+--------+
|      6 |
|      7 |
|      8 |
+--------+
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.
- For more information about InnoDB FULLTEXT search, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and Section 12.9, “Full-Text Search Functions”.

24.31.12 The INFORMATION_SCHEMA INNODB_FT_INDEX_CACHE Table

INNODB_FT_INDEX_CACHE: Contains token information about newly inserted rows in a FULLTEXT index. To avoid expensive index reorganization during DML operations, the information about newly indexed words is stored separately, and combined with the main search index only when OPTIMIZE TABLE is run, when the server is shut down, or when the cache size exceeds a limit defined by innodb_ft_cache_size or innodb_ft_total_cache_size.

Before you query this table, set the configuration variable innodb_ft_aux_table to the name (including the database name) of the table that contains the FULLTEXT index, for example test/articles.

For related usage information and examples, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

Table 24.12 INNODB_FT_INDEX_CACHE Columns

Column name	Description
WORD	A word extracted from the text of a newly inserted row.
FIRST_DOC_ID	The first document ID that this word appears in the FULLTEXT index.
LAST_DOC_ID	The last document ID that this word appears in the FULLTEXT index.
DOC_COUNT	The number of rows this word appears in the FULLTEXT index. The same word can occur several times within the cache table, once for each combination of DOC_ID and POSITION values.

Column name	Description
DOC_ID	The document ID of the newly inserted row. This value might reflect the value of an ID column that you defined for the underlying table, or it can be a sequence value generated by InnoDB when the table does not contain a suitable column.
POSITION	The position of this particular instance of the word within the relevant document identified by the DOC_ID value. The value does not represent an absolute position; it is an offset added to the POSITION of the previous instance of that word.

Notes:

- This table initially appears empty, until you set the value of the configuration variable innodb_ft_aux_table. The following example demonstrates how to use the innodb_ft_aux_table option to show information about a FULLTEXT index for a specified table.

```
mysql> USE test;

mysql> CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT,
      FULLTEXT (title,body)
    ) ENGINE=InnoDB;

mysql> INSERT INTO articles (title,body) VALUES
    ('MySQL Tutorial','DBMS stands for DataBase ...'),
    ('How To Use MySQL Well','After you went through a ...'),
    ('Optimizing MySQL','In this tutorial we will show ...'),
    ('1001 MySQL Tricks','1. Never run mysqld as root. 2. ...'),
    ('MySQL vs. YourSQL','In the following database comparison ...'),
    ('MySQL Security','When configured properly, MySQL ...');

mysql> SET GLOBAL innodb_ft_aux_table = 'test/articles';
Query OK, 0 rows affected (0.00 sec)

mysql> USE INFORMATION_SCHEMA;

mysql> SELECT word, doc_count, doc_id, position FROM INNODB_FT_INDEX_CACHE LIMIT 5;
+-----------+-----------+--------+----------+
| word      | doc_count | doc_id | position |
+-----------+-----------+--------+----------+
| 1001      |         1 |      4 |        0 |
| after     |         1 |      2 |       22 |
| comparison |         1 |      5 |       44 |
| configured |         1 |      6 |       20 |
| database  |         2 |      1 |       31 |
+-----------+-----------+--------+----------+
```

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.
- For more information about InnoDB FULLTEXT search, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and Section 12.9, “Full-Text Search Functions”.

24.31.13 The INFORMATION_SCHEMA INNODB_FT_INDEX_TABLE Table

The INNODB_FT_INDEX_TABLE table displays information about the inverted index used to process text searches against the FULLTEXT index of an InnoDB table.

For related usage information and examples, see Section 15.14.4, “InnoDB INFORMATION_SCHEMA FULLTEXT Index Tables”.

Before you query this table, set the configuration variable `innodb_ft_aux_table` to the name (including the database name) of the table that contains the `FULLTEXT` index, for example `test/articles`.

Table 24.13 INNODB_FT_INDEX_TABLE Columns

Column name	Description
`WORD`	A word extracted from the text of the columns that are part of a `FULLTEXT`.
`FIRST_DOC_ID`	The first document ID that this word appears in the `FULLTEXT` index.
`LAST_DOC_ID`	The last document ID that this word appears in the `FULLTEXT` index.
`DOC_COUNT`	The number of rows this word appears in the `FULLTEXT` index. The same word can occur several times within the cache table, once for each combination of `DOC_ID` and `POSITION` values.
`DOC_ID`	The document ID of the row containing the word. This value might reflect the value of an ID column that you defined for the underlying table, or it can be a sequence value generated by `InnoDB` when the table does not contain a suitable column.
`POSITION`	The position of this particular instance of the word within the relevant document identified by the `DOC_ID` value.

Notes:

- This table initially appears empty, until you set the value of the configuration variable `innodb_ft_aux_table`. The following example demonstrates how to use the `innodb_ft_aux_table` option to show information about a `FULLTEXT` index for a specified table. Before information for newly inserted rows appears in `INNODB_FT_INDEX_TABLE`, the `FULLTEXT` index cache must be flushed to disk. This is accomplished by running an `OPTIMIZE TABLE` operation on the indexed table with `innodb_optimize_fulltext_only=ON`.

```
mysql> use test;

mysql> CREATE TABLE articles (
      id INT UNSIGNED AUTO_INCREMENT NOT NULL PRIMARY KEY,
      title VARCHAR(200),
      body TEXT,
      FULLTEXT (title,body)
    ) ENGINE=InnoDB;

mysql> INSERT INTO articles (title,body) VALUES
    ('MySQL Tutorial','DBMS stands for DataBase ...'),
    ('How To Use MySQL Well','After you went through a ...'),
    ('Optimizing MySQL','In this tutorial we will show ...'),
    ('1001 MySQL Tricks','1. Never run mysqld as root. 2. ...'),
    ('MySQL vs. YourSQL','In the following database comparison ...'),
    ('MySQL Security','When configured properly, MySQL ...');

mysql> SET GLOBAL innodb_optimize_fulltext_only=ON;
Query OK, 0 rows affected (0.00 sec)

mysql> OPTIMIZE TABLE articles;
+---------------+----------+----------+----------+
| Table         | Op       | Msg_type | Msg_text |
+---------------+----------+----------+----------+
| test.articles | optimize | status   | OK       |
+---------------+----------+----------+----------+

mysql> SET GLOBAL innodb_ft_aux_table = 'test/articles';
Query OK, 0 rows affected (0.00 sec)

mysql> USE INFORMATION_SCHEMA;
```

```
mysql> SELECT word, doc_count, doc_id, position FROM INNODB_FT_INDEX_TABLE LIMIT 5;
+------------+-----------+--------+----------+
| word       | doc_count | doc_id | position |
+------------+-----------+--------+----------+
| 1001       |         1 |      4 |        0 |
| after      |         1 |      2 |       22 |
| comparison |         1 |      5 |       44 |
| configured |         1 |      6 |       20 |
| database   |         2 |      1 |       31 |
+------------+-----------+--------+----------+
```

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For more information about `InnoDB FULLTEXT` search, see Section 15.8.2.4, “InnoDB FULLTEXT Indexes”, and Section 12.9, “Full-Text Search Functions”.

24.31.14 The INFORMATION_SCHEMA INNODB_LOCKS Table

The `INNODB_LOCKS` table contains information about each lock that an `InnoDB` transaction has requested but not yet acquired, and each lock that a transaction holds that is blocking another transaction.

Note

This table is deprecated and is removed as of MySQL 8.0.1. Use the Performance Schema `data_locks` table instead. See Section 25.11.12.1, “The data_locks Table”.

Differences between `INNODB_LOCKS` and `data_locks`:

- If a transaction holds a lock, `INNODB_LOCKS` displays the lock only if another transaction is waiting for it. `data_locks` displays the lock regardless of whether any transaction is waiting for it.
- The `data_locks` table has no columns corresponding to `LOCK_SPACE`, `LOCK_PAGE`, or `LOCK_REC`.
- The `INNODB_LOCKS` table requires the global `PROCESS` privilege. The `data_locks` table requires the usual Performance Schema privilege of `SELECT` on the table to be selected from.

The following table shows the mapping from `INNODB_LOCKS` columns to `data_locks` columns. Use this information to migrate applications from one table to the other.

Table 24.14 Mapping from INNODB_LOCKS to data_locks Columns

INNODB_LOCKS Column	data_locks Column
`LOCK_ID`	`ENGINE_LOCK_ID`
`LOCK_TRX_ID`	`ENGINE_TRANSACTION_ID`
`LOCK_MODE`	`LOCK_MODE`
`LOCK_TYPE`	`LOCK_TYPE`
`LOCK_TABLE` (combined schema/table names)	`OBJECT_SCHEMA` (schema name), `OBJECT_NAME` (table name)
`LOCK_INDEX`	`INDEX_NAME`
`LOCK_SPACE`	None
`LOCK_PAGE`	None

INNODB_LOCKS Column	data_locks Column
LOCK_REC	None
LOCK_DATA	LOCK_DATA

24.31.15 The INFORMATION_SCHEMA INNODB_LOCK_WAITS Table

The INNODB_LOCK_WAITS table contains one or more rows for each blocked InnoDB transaction, indicating the lock it has requested and any locks that are blocking that request.

Note

This table is deprecated and is removed as of MySQL 8.0.1. Use the Performance Schema data_lock_waits table instead. See Section 25.11.12.2, "The data_lock_waits Table".

The tables differ in the privileges required: The INNODB_LOCK_WAITS table requires the global PROCESS privilege. The data_lock_waits table requires the usual Performance Schema privilege of SELECT on the table to be selected from.

The following table shows the mapping from INNODB_LOCK_WAITS columns to data_lock_waits columns. Use this information to migrate applications from one table to the other.

Table 24.15 Mapping from INNODB_LOCK_WAITS to data_lock_waits Columns

INNODB_LOCK_WAITS Column	data_lock_waits Column
REQUESTING_TRX_ID	REQUESTING_ENGINE_TRANSACTION_ID
REQUESTED_LOCK_ID	REQUESTING_ENGINE_LOCK_ID
BLOCKING_TRX_ID	BLOCKING_ENGINE_TRANSACTION_ID
BLOCKING_LOCK_ID	BLOCKING_ENGINE_LOCK_ID

24.31.16 The INFORMATION_SCHEMA INNODB_METRICS Table

This INFORMATION_SCHEMA table presents a wide variety of InnoDB performance information, complementing the specific focus areas of the PERFORMANCE_SCHEMA tables for InnoDB. With simple queries, you can check the overall health of the system. With more detailed queries, you can diagnose issues such as performance bottlenecks, resource shortages, and application issues.

Each monitor represents a point within the InnoDB source code that is instrumented to gather counter information. Each counter can be started, stopped, and reset. You can also perform these actions for a group of counters using their common module name.

By default, relatively little data is collected. To start, stop, and reset counters, you set one of the configuration options innodb_monitor_enable, innodb_monitor_disable, innodb_monitor_reset, or innodb_monitor_reset_all, using the name of the counter, the name of the module, a wildcard match for such a name using the "%" character, or the special keyword all.

For usage information, see Section 15.14.6, "InnoDB INFORMATION_SCHEMA Metrics Table".

Table 24.16 INNODB_METRICS Columns

Column name	Description
NAME	Unique name for the counter.
SUBSYSTEM	The aspect of InnoDB that the metric applies to. See the list following the table for the corresponding module names to use with the SET GLOBAL syntax.

Column name	Description
COUNT	Value since the counter is enabled.
MAX_COUNT	Maximum value since the counter is enabled.
MIN_COUNT	Minimum value since the counter is enabled.
AVG_COUNT	Average value since the counter is enabled.
COUNT_RESET	Counter value since it was last reset. (The _RESET fields act like the lap counter on a stopwatch: you can measure the activity during some time interval, while the cumulative figures are still available in the COUNT, MAX_COUNT, and so on fields.)
MAX_COUNT_RESET	Maximum counter value since it was last reset.
MIN_COUNT_RESET	Minimum counter value since it was last reset.
AVG_COUNT_RESET	Average counter value since it was last reset.
TIME_ENABLED	Timestamp of last start.
TIME_DISABLED	Timestamp of last stop.
TIME_ELAPSED	Elapsed time in seconds since the counter started.
TIME_RESET	Timestamp of last stop.
STATUS	Whether the counter is still running (enabled) or stopped (disabled).
TYPE	Whether the item is a cumulative counter, or measures the current value of some resource.
COMMENT	Counter description.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_METRICS WHERE NAME="dml_inserts"\G
*************************** 1. row ***************************
           NAME: dml_inserts
      SUBSYSTEM: dml
          COUNT: 3
      MAX_COUNT: 3
      MIN_COUNT: NULL
      AVG_COUNT: 0.046153846153846156
    COUNT_RESET: 3
MAX_COUNT_RESET: 3
MIN_COUNT_RESET: NULL
AVG_COUNT_RESET: NULL
   TIME_ENABLED: 2014-12-04 14:18:28
  TIME_DISABLED: NULL
   TIME_ELAPSED: 65
     TIME_RESET: NULL
         STATUS: enabled
           TYPE: status_counter
        COMMENT: Number of rows inserted
```

Notes:

- You must have the PROCESS privilege to query this table.
- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.

24.31.17 The INFORMATION_SCHEMA INNODB_SYS_COLUMNS Table

The INNODB_SYS_COLUMNS table provides metadata about InnoDB table columns.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Table 24.17 INNODB_SYS_COLUMNS Columns

Column name	Description
TABLE_ID	An identifier representing the table associated with the column; the same value from INNODB_SYS_TABLES.TABLE_ID.
NAME	The name of each column in each table. These names can be uppercase or lowercase depending on the lower_case_table_names setting. There are no special system-reserved names for columns.
POS	The ordinal position of the column within the table, starting from 0 and incrementing sequentially. When a column is dropped, the remaining columns are reordered so that the sequence has no gaps. The POS value for a virtual generated column encodes the column sequence number and ordinal position of the column. For more information, see the POS column description in Section 24.31.26, "The INFORMATION_SCHEMA INNODB_SYS_VIRTUAL Table".
MTYPE	Stands for "main type". A numeric identifier for the column type. 1 = VARCHAR, 2 = CHAR, 3 = FIXBINARY, 4 = BINARY, 5 = BLOB, 6 = INT, 7 = SYS_CHILD, 8 = SYS, 9 = FLOAT, 10 = DOUBLE, 11 = DECIMAL, 12 = VARMYSQL, 13 = MYSQL, 14 = GEOMETRY.
PRTYPE	The InnoDB "precise type", a binary value with bits representing MySQL data type, character set code, and nullability.
LEN	The column length, for example 4 for INT and 8 for BIGINT. For character columns in multibyte character sets, this length value is the maximum length in bytes needed to represent a definition such as VARCHAR(*N*); that is, it might be 2**N*, 3**N*, and so on depending on the character encoding.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_COLUMNS where TABLE_ID = 71\G
*************************** 1. row ***************************
TABLE_ID: 71
    NAME: col1
     POS: 0
   MTYPE: 6
  PRTYPE: 1027
     LEN: 4
*************************** 2. row ***************************
TABLE_ID: 71
    NAME: col2
     POS: 1
   MTYPE: 2
  PRTYPE: 524542
     LEN: 10
*************************** 3. row ***************************
TABLE_ID: 71
    NAME: col3
     POS: 2
   MTYPE: 1
  PRTYPE: 524303
     LEN: 10
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

24.31.18 The INFORMATION_SCHEMA INNODB_SYS_DATAFILES Table

The INNODB_SYS_DATAFILES table provides data file path information for InnoDB file-per-table and general tablespaces.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Note

The INFORMATION_SCHEMA.FILES table provides data file path information and other metadata about all InnoDB tablespace types including file-per-table tablespaces, general tablespaces, the system tablespace, temporary tablespaces, and undo tablespaces (if present).

Table 24.18 INNODB_SYS_DATAFILES Columns

Column name	Description
SPACE	The tablespace Space ID.
PATH	The tablespace data file path (for example, .\world\innodb\city.ibd). If a file-per-table tablespace is created in a location outside the MySQL data directory using the DATA DIRECTORY clause of the CREATE TABLE statement, the tablespace PATH field shows the fully qualified directory path.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_DATAFILES WHERE SPACE = 57\G
*************************** 1. row ***************************
SPACE: 57
 PATH: ./test/t1.ibd
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

24.31.19 The INFORMATION_SCHEMA INNODB_SYS_FIELDS Table

The INNODB_SYS_FIELDS table provides metadata about the key columns (fields) of InnoDB indexes.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Table 24.19 INNODB_SYS_FIELDS Columns

Column name	Description
INDEX_ID	An identifier for the index associated with this key field, using the same value as in INNODB_SYS_INDEXES.INDEX_ID.
NAME	The name of the original column from the table, using the same value as in INNODB_SYS_COLUMNS.NAME.
POS	The ordinal position of the key field within the index, starting from 0 and incrementing sequentially. When a column is dropped, the remaining columns are reordered so that the sequence has no gaps.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FIELDS WHERE INDEX_ID = 117\G
*************************** 1. row ***************************
INDEX_ID: 117
    NAME: col1
     POS: 0
```

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.

24.31.20 The INFORMATION_SCHEMA INNODB_SYS_FOREIGN Table

The `INNODB_SYS_FOREIGN` table provides metadata about `InnoDB` foreign keys.

For related usage information and examples, see Section 15.14.3, "InnoDB INFORMATION_SCHEMA System Tables".

Table 24.20 INNODB_SYS_FOREIGN Columns

Column name	Description
`ID`	The name (not a numeric value) of the foreign key index. Preceded by the database name, for example, `test/products_fk`.
`FOR_NAME`	The name of the child table in this foreign key relationship.
`REF_NAME`	The name of the parent table in this foreign key relationship.
`N_COLS`	The number of columns in the foreign key index.
`TYPE`	A collection of bit flags with information about the foreign key column, ORed together. 1 = `ON DELETE CASCADE`, 2 = `ON UPDATE SET` NULL, 4 = `ON UPDATE CASCADE`, 8 = `ON UPDATE SET` NULL, 16 = `ON DELETE NO` ACTION, 32 = `ON UPDATE NO` ACTION.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FOREIGN\G
*************************** 1. row ***************************
      ID: test/fk1
FOR_NAME: test/child
REF_NAME: test/parent
  N_COLS: 1
    TYPE: 1
```

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.

24.31.21 The INFORMATION_SCHEMA INNODB_SYS_FOREIGN_COLS Table

The `INNODB_SYS_FOREIGN_COLS` table provides status information about the columns of `InnoDB` foreign keys.

For related usage information and examples, see Section 15.14.3, "InnoDB INFORMATION_SCHEMA System Tables".

Table 24.21 INNODB_SYS_FOREIGN_COLS Columns

Column name	Description
ID	The foreign key index associated with this index key field, using the same value as INNODB_SYS_FOREIGN.ID.
FOR_COL_NAME	The name of the associated column in the child table.
REF_COL_NAME	The name of the associated column in the parent table.
POS	The ordinal position of this key field within the foreign key index, starting from 0.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_FOREIGN_COLS WHERE ID = 'test/fk1'\G
*************************** 1. row ***************************
          ID: test/fk1
FOR_COL_NAME: parent_id
REF_COL_NAME: id
         POS: 0
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

24.31.22 The INFORMATION_SCHEMA INNODB_SYS_INDEXES Table

The INNODB_SYS_INDEXES table provides metadata about InnoDB indexes.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Table 24.22 INNODB_SYS_INDEXES Columns

Column name	Description
INDEX_ID	An identifier for each index that is unique across all the databases in an instance.
NAME	The name of the index. Most indexes created implicitly by InnoDB have consistent names but the index names are not necessarily unique. For example, PRIMARY for a primary key index, GEN_CLUST_INDEX for the index representing a primary key when one is not specified, and ID_IND, FOR_IND, and REF_IND for foreign key constraints.
TABLE_ID	An identifier representing the table associated with the index; the same value from INNODB_SYS_TABLES.TABLE_ID.
TYPE	A numeric identifier signifying the kind of index. 0 = Secondary Index, 1 = Clustered Index, 2 = Unique Index, 3 = Primary Index, 32 = Full-text Index, 64 = Spatial Index, 128 = A secondary index that includes a virtual generated column.
N_FIELDS	The number of columns in the index key. For the GEN_CLUST_INDEX indexes, this value is 0 because the index is created using an artificial value rather than a real table column.
PAGE_NO	The root page number of the index B-tree. For full-text indexes, the PAGE_NO field is unused and set to -1 (FIL_NULL) because the full-text index is laid out in several B-trees (auxiliary tables).

Column name	Description
SPACE	An identifier for the tablespace where the index resides. 0 means the InnoDB system tablespace. Any other number represents a table created in file-per-table mode with a separate .ibd file. This identifier stays the same after a TRUNCATE TABLE statement. Because all indexes for a table reside in the same tablespace as the table, this value is not necessarily unique.
MERGE_THRESHOLD	The merge threshold value for index pages. If the amount of data in an index page falls below the MERGE_THRESHOLD value when a row is deleted or when a row is shortened by an update operation, InnoDB attempts to merge the index page with the neighboring index page. The default threshold value is 50%. For more information, see Section 15.6.12, “Configuring the Merge Threshold for Index Pages”.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_INDEXES WHERE TABLE_ID = 34\G
*************************** 1. row ***************************
       INDEX_ID: 39
           NAME: GEN_CLUST_INDEX
       TABLE_ID: 34
           TYPE: 1
       N_FIELDS: 0
        PAGE_NO: 3
          SPACE: 23
MERGE_THRESHOLD: 50
*************************** 2. row ***************************
       INDEX_ID: 40
           NAME: i1
       TABLE_ID: 34
           TYPE: 0
       N_FIELDS: 1
        PAGE_NO: 4
          SPACE: 23
MERGE_THRESHOLD: 50
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

24.31.23 The INFORMATION_SCHEMA INNODB_SYS_TABLES Table

The INNODB_SYS_TABLES table provides metadata about InnoDB tables.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Table 24.23 INNODB_SYS_TABLES Columns

Column name	Description
TABLE_ID	An identifier for each InnoDB table that is unique across all databases in the instance.
NAME	The name of the table. Preceded by the database name where appropriate, for example test/t1. InnoDB system table names are in all uppercase. Names of databases and user tables are in the same case as they were originally defined, possibly influenced by the lower_case_table_names setting.

Column name	Description
FLAG	This value provides bit level information about table format and storage characteristics including row format, compressed page size (if applicable), and whether or not the DATA DIRECTORY clause was used with CREATE TABLE or ALTER TABLE.
N_COLS	The number of columns in the table. The number reported includes three hidden columns that are created by InnoDB (DB_ROW_ID, DB_TRX_ID, and DB_ROLL_PTR). The number reported also includes virtual generated columns, if present.
SPACE	An identifier for the tablespace where the table resides. 0 means the InnoDB system tablespace. Any other number represents either a file-per-table tablespace or a general tablespace. This identifier stays the same after a TRUNCATE TABLE statement. For file-per-table tablespaces, this identifier is unique for tables across all databases in the instance.
ROW_FORMAT	The table's row format (Compact, Redundant, Dynamic, or Compressed).
ZIP_PAGE_SIZE	The zip page size. Only applies to tables that use the Compressed row format.
SPACE_TYPE	The type of tablespace to which the table belongs. Possible values include System (for the InnoDB system tablespace), General (for InnoDB general tablespaces created using CREATE TABLESPACE, and Single (for InnoDB file-per-table tablespaces). Tables assigned to the system tablespace using the CREATE TABLE or ALTER TABLE TABLESPACE=innodb_system clause have a General SPACE_TYPE.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE TABLE_ID = 214\G
*************************** 1. row ***************************
     TABLE_ID: 214
         NAME: test/t1
         FLAG: 129
       N_COLS: 4
        SPACE: 233
   ROW_FORMAT: Compact
ZIP_PAGE_SIZE: 0
   SPACE_TYPE: General
```

Notes:

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

Interpreting the INNODB_SYS_TABLES.FLAG Column Value:

The INNODB_SYS_TABLES.FLAG column provides bit-level information about the table's format and storage characteristics. You can interpret the FLAG column value by adding together the applicable decimal numeric values that are provided in the following table.

Table 24.24 Bit Position Values for Interpreting INNODB_SYS_TABLES FLAG Column Data

Bit Position	Description	Decimal Numeric Value
0	This bit is set if the row format is not REDUNDANT. In other words, it is set if	• 0 - REDUNDANT • 1 - COMPACT, DYNAMIC or COMPRESSED

Bit Position	Description	Decimal Numeric Value
	the row format is `COMPACT`, `DYNAMIC` or `COMPRESSED`.	
1-4	These four bits contain a small number that represents the compressed page size of the table. The `INNODB_SYS_TABLES.ZIP_PAGE_SIZE` field also reports the compressed page size, if applicable.	• 0 - Not Compressed • 2 - 1024 Byte Compressed Page Size • 4 - 2048 Byte Compressed Page Size • 6 - 4096 Byte Compressed Page Size • 8 - 8192 Byte Compressed Page Size • 10 - 16384 Byte Compressed Page Size
5	This bit is set if the row format is `DYNAMIC` or `COMPRESSED`.	• 0 - `REDUNDANT` or `COMPACT` • 32 - `DYNAMIC` or `COMPRESSED`
6	This bit is set if the `DATA DIRECTORY` option is used with `CREATE TABLE` or `ALTER TABLE`. This bit is set for file-per-table tablespaces that are located in directories other than the default data directory (`datadir`).	• 0 - Not a remote file-per-table tablespace • 64 - A remote file-per-table tablespace
7	This bit is set if the table is assigned to a shared tablespace (either a general tablespace or a system tablespace) using the `CREATE TABLE` or `ALTER TABLE TABLESPACE=`*`tablespace_name`* option.	• 0 - Table is located in a default location depending on the value of the `innodb_file_per_table`. • 128 - The table is explicitly assigned to a shared tablespace.

In the following, table `t1` uses `ROW_FORMAT=DYNAMIC` and has a `FLAG` value of 33. Based on the information in the preceding table, we can see that bit position 0 would be set to 1, and bit position 5 would be set to 32 for a table with a `DYNAMIC` row format. These values add up to a `FLAG` value of 33.

```
mysql> use test;
Database changed

mysql> CREATE TABLE t1 (c1 int) ROW_FORMAT=DYNAMIC;
Query OK, 0 rows affected (0.02 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME LIKE 'test/t1'\G
*************************** 1. row ***************************
     TABLE_ID: 89
         NAME: test/t1
         FLAG: 33
       N_COLS: 4
        SPACE: 75
   ROW_FORMAT: Dynamic
ZIP_PAGE_SIZE: 0
1 row in set (0.00 sec)
```

24.31.24 The INFORMATION_SCHEMA INNODB_SYS_TABLESPACES Table

The `INNODB_SYS_TABLESPACES` table provides metadata about `InnoDB` file-per-table and general tablespaces.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Note

The `INFORMATION_SCHEMA.FILES` table provides metadata about all `InnoDB` tablespace types including file-per-table tablespaces, general tablespaces, the system tablespace, temporary tablespaces, and undo tablespaces (if present).

Table 24.25 INNODB_SYS_TABLESPACES Columns

Column name	Description
`SPACE`	Tablespace Space ID.
`NAME`	The database and table name (for example, `world_innodb\city`)
`FLAG`	This value provides bit level information about tablespace format and storage characteristics.
`ROW_FORMAT`	The tablespace row format (Compact or Redundant, Dynamic, or Compressed). The data in this field is interpreted from the tablespace flags information that resides in the .ibd file.
`PAGE_SIZE`	The tablespace page size. The data in this field is interpreted from the tablespace flags information that resides in the .ibd file.
`ZIP_PAGE_SIZE`	The tablespace zip page size. The data in this field is interpreted from the tablespace flags information that resides in the .ibd file.
`SPACE_TYPE`	The type of tablespace. Possible values include `General` (for `InnoDB` general tablespaces created using `CREATE TABLESPACE` and `Single` (for `InnoDB` file-per-table tablespaces).
`FS_BLOCK_SIZE`	The file system block size, which is the unit size used for hole punching. This column was added with the introduction of the `InnoDB` transparent page compression feature.
`FILE_SIZE`	The apparent size of the file, which represents the maximum size of the file, uncompressed. This column was added with the introduction of the `InnoDB` transparent page compression feature.
`ALLOCATED_SIZE`	The actual size of the file, which is the amount of space allocated on disk. This column was added with the introduction of the `InnoDB` transparent page compression feature.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES WHERE SPACE = 26\G
*************************** 1. row ***************************
         SPACE: 26
          NAME: test/t1
          FLAG: 0
    ROW_FORMAT: Compact or Redundant
     PAGE_SIZE: 16384
 ZIP_PAGE_SIZE: 0
    SPACE_TYPE: Single
 FS_BLOCK_SIZE: 4096
     FILE_SIZE: 98304
ALLOCATED_SIZE: 65536
```

Notes:

- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- There is no way to determine from this flag integer if the tablespace row format is Redundant or Compact. As a result, the possible values for the `ROW_FORMAT` field are “Compact or Redundant”, “Compressed”, or “Dynamic.”

Interpreting the INNODB_SYS_TABLESPACES.FLAG Column Value:

The `INNODB_SYS_TABLESPACES.FLAG` column provides bit-level information about tablespace format and storage characteristics.

You can interpret the tablespace `FLAG` column value by adding together the applicable decimal numeric values that are provided in the following table.

Table 24.26 Bit Position Values for Interpreting INNODB_SYS_TABLESPACES FLAG Column Data

Bit Position	Description	Decimal Numeric Value
0	This bit is set if the row format of tables in the tablespace is `DYNAMIC` or `COMPRESSED`. If the bit is not set, the row format of tables in the tablespace may be either `REDUNDANT` or `COMPACT`. If it is a file-per-table tablespace, you can query `INNODB_SYS_TABLES` to determine if the row format is `REDUNDANT` or `COMPACT`.	• 0 - `REDUNDANT` or `COMPACT` • 1 - `DYNAMIC` or `COMPRESSED`
1-4	These four bits contain a small number that represents the compressed page size (the `KEY_BLOCK_SIZE` or “physical block size”) of the tablespace.	• 0 - Not Compressed • 2 - 1024 Byte Compressed Page Size • 4 - 2048 Byte Compressed Page Size • 6 - 4096 Byte Compressed Page Size • 8 - 8192 Byte Compressed Page Size • 10 - 16384 Byte Compressed Page Size • 12 - 32768 Byte Compressed Page Size • 14 - 65536 Byte Compressed Page Size
5	This bit is set for file-per-table tablespaces if the row format of the table is `DYNAMIC` or `COMPRESSED`. General tablespaces that do not contain compressed tables will have the first 6 bits set to zero, including this bit, making it appear that the tablespace holds `REDUNDANT` or `COMPACT` tables. But actually, general tablespaces may contain any combination of `REDUNDANT`, `COMPACT` and `DYNAMIC` tables. For more information about general tablespaces, see `CREATE TABLESPACE`.	• 0 - `REDUNDANT` or `COMPACT` • 32 - `DYNAMIC` or `COMPRESSED`
6-9	These four bits contain a small number that represents the uncompressed page size (logical page size) of the tablespace. The setting is zero if the logical page size is the original `InnoDB` default page size of 16K.	• 192 - 4096 Byte Logical/ Uncompressed Page Size • 256 - 8192 Byte Logical/ Uncompressed Page Size • 0 - 16384 Byte Logical/ Uncompressed Page size

Bit Position	Description	Decimal Numeric Value
		• 384 - 32768 Byte Logical/ Uncompressed Page Size • 448 - 65536 Byte Logical/ Uncompressed Page Size
10	This bit is set if the `DATA DIRECTORY` option is used with `CREATE TABLE` or `ALTER TABLE`. This bit is set for file-per-table tablespaces that are located in directories other than the default data directory (`datadir`).	• 0 - Not a remote file-per-table tablespace • 1024 - A remote file-per-table tablespace
11	This bit is set if the tablespace is a shared general tablespace created using `CREATE TABLESPACE`.	• 0 - Table is located in a default location depending on the value of the `innodb_file_per_table` setting. • 2048 - The table was explicitly assigned to a shared tablespace.
12	This bit is set if the tablespace is dedicated to temporary tables. Only the predefined `ibtmp1` tablespace uses this flag.	• 0 - The tablespace does not contain temporary tables, so it is not recreated upon startup. • 4096 - The tablespace contains temporary tables and is recreated on startup.

In the following example, table `t1` is created with `innodb_file_per_table=ON`, which creates table `t1` in its own tablespace. When querying `INNODB_SYS_TABLESPACES`, we see that the tablespace has a `FLAG` value of 33. To determine how this value is arrived at, review the bit values described in the preceding table. Bit 0 has a value of 1 because table `t1` uses the `DYNAMIC` row format. Bit 5 has a value of 32 because the tablespace is a file-per-table tablespace that uses a `DYNAMIC` row format. Bit position 6-9 is 0 because `innodb_page_size` is set to the default 16K value. The other bit values are not applicable and are therefore set to 0. The values for bit position 0 and bit position 5 add up to a `FLAG` value of 33.

```
mysql> use test;
Database changed

mysql> SHOW VARIABLES LIKE 'innodb_file_per_table';
+-----------------------+-------+
| Variable_name         | Value |
+-----------------------+-------+
| innodb_file_per_table | ON    |
+-----------------------+-------+

mysql> SHOW VARIABLES LIKE 'innodb_page_size';
+------------------+-------+
| Variable_name    | Value |
+------------------+-------+
| innodb_page_size | 16384 |
+------------------+-------+

mysql> CREATE TABLE t1 (c1 int) ROW_FORMAT=DYNAMIC;
Query OK, 0 rows affected (0.02 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLESPACES WHERE NAME LIKE 'test/t1'\G
*************************** 1. row ***************************
        SPACE: 75
```

```
         NAME: test/t1
         FLAG: 33
   ROW_FORMAT: Dynamic
    PAGE_SIZE: 16384
ZIP_PAGE_SIZE: 0
1 row in set (0.00 sec)
```

24.31.25 The INFORMATION_SCHEMA INNODB_SYS_TABLESTATS View

The `INNODB_SYS_TABLESTATS` provides a view of low-level status information about `InnoDB` tables. This data is used by the MySQL optimizer to calculate which index to use when querying an `InnoDB` table. This information is derived from in-memory data structures rather than corresponding to data stored on disk. There is no corresponding internal `InnoDB` system table.

`InnoDB` tables are represented in this view if they have been opened since the last server restart, and not aged out of the table cache. Tables for which persistent stats are available are always represented in this view.

Table statistics are only updated for `DELETE` or `UPDATE` operations that modify indexed columns. Statistics are not updated by operations that only modify non-indexed columns.

For related usage information and examples, see Section 15.14.3, “InnoDB INFORMATION_SCHEMA System Tables”.

Table 24.27 INNODB_SYS_TABLESTATS Columns

Column name	Description
`TABLE_ID`	An identifier representing the table for which statistics are available, using the same value as `INNODB_SYS_TABLES.TABLE_ID`.
`NAME`	The name of the table, using the same value as `INNODB_SYS_TABLES.NAME`.
`STATS_INITIALIZED`	The value is `Initialized` if the statistics are already collected, `Uninitialized` if not.
`NUM_ROWS`	The current estimated number of rows in the table. Updated after each DML operation. Could be imprecise if uncommitted transactions are inserting into or deleting from the table.
`CLUST_INDEX_SIZE`	Number of pages on disk that store the clustered index, which holds the `InnoDB` table data in primary key order. This value might be null if no statistics are collected yet for the table.
`OTHER_INDEX_SIZE`	Number of pages on disk that store all secondary indexes for the table. This value might be null if no statistics are collected yet for the table.
`MODIFIED_COUNTER`	The number of rows modified by DML operations, such as `INSERT`, `UPDATE`, `DELETE`, and also cascade operations from foreign keys. This column is reset each time table statistics are recalculated
`AUTOINC`	The next number to be issued for any auto-increment-based operation. The rate at which the `AUTOINC` value changes depends on how many times auto-increment numbers have been requested and how many numbers are granted per request.
`REF_COUNT`	When this counter reaches zero, the table metadata can be evicted from the table cache.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_TABLESTATS where TABLE_ID = 71\G
*************************** 1. row ***************************
         TABLE_ID: 71
             NAME: test/t1
STATS_INITIALIZED: Initialized
```

```
        NUM_ROWS: 1
 CLUST_INDEX_SIZE: 1
 OTHER_INDEX_SIZE: 0
 MODIFIED_COUNTER: 1
          AUTOINC: 0
        REF_COUNT: 1
```

Notes:

- This table is primarily useful for expert-level performance monitoring, or when developing performance-related extensions for MySQL.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.

24.31.26 The INFORMATION_SCHEMA INNODB_SYS_VIRTUAL Table

The `INNODB_SYS_VIRTUAL` table provides metadata about `InnoDB` virtual generated columns and columns upon which virtual generated columns are based.

A row appears in the `INNODB_SYS_VIRTUAL` table for each column upon which a virtual generated column is based.

Table 24.28 INNODB_SYS_VIRTUAL Columns

Column name	Description
`TABLE_ID`	An identifier representing the table associated with the virtual column; the same value as `INNODB_SYS_TABLES.TABLE_ID`.
`POS`	The position value of the virtual generated column. The value is large because it encodes the column sequence number and ordinal position. The formula used to calculate the value uses a bitwise operation. The formula is `((nth virtual generated column for the InnoDB instance + 1) << 16) + the ordinal position of the virtual generated column`. For example, if the first virtual generated column in the `InnoDB` instance is the third column of the table, the formula is (0 + 1) << 16) + 2. The first virtual generated column in the `InnoDB` instance is always number 0. As the third column in the table, the ordinal position of the virtual generated column is 2. Ordinal positions are counted from 0.
`BASE_POS`	The ordinal position of the columns upon which a virtual generated column is based.

Example:

```
mysql> CREATE TABLE `t1` (
         `a` int(11) DEFAULT NULL,
         `b` int(11) DEFAULT NULL,
         `c` int(11) GENERATED ALWAYS AS (a+b) VIRTUAL,
         `h` varchar(10) DEFAULT NULL
       ) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4;
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_SYS_VIRTUAL
       WHERE TABLE_ID IN (SELECT TABLE_ID FROM INFORMATION_SCHEMA.INNODB_SYS_TABLES WHERE NAME LIKE "test/t
+----------+-------+----------+
| TABLE_ID | POS   | BASE_POS |
+----------+-------+----------+
|       98 | 65538 |        0 |
|       98 | 65538 |        1 |
+----------+-------+----------+
```

Notes:

- If a constant value is assigned to a virtual generated column, as in the following example, an entry for the column does not appear in the INNODB_SYS_VIRTUAL table. For an entry to appear, a virtual generated column must have a base column.

```
mysql> CREATE TABLE `t1` (
         `a` int(11) DEFAULT NULL,
         `b` int(11) DEFAULT NULL,
         `c` int(11) GENERATED ALWAYS AS (5) VIRTUAL
       ) ENGINE=InnoDB DEFAULT CHARSET=utf8mb4;
```

  However, metadata for such a column appears in the INNODB_SYS_COLUMNS table.

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.

- You must have the PROCESS privilege to query this table.

24.31.27 The INFORMATION_SCHEMA INNODB_TEMP_TABLE_INFO Table

INNODB_TEMP_TABLE_INFO contains metadata about active InnoDB temporary tables. With the exception of optimized internal temporary tables used by InnoDB, INNODB_TEMP_TABLE_INFO reports on all user and system-created temporary tables that are active within a given InnoDB instance. The table is maintained in memory and not persisted to disk.

For usage information and examples, see Section 15.14.7, “InnoDB INFORMATION_SCHEMA Temporary Table Information Table”.

Table 24.29 INNODB_TEMP_TABLE_INFO Columns

Column name	Description
TABLE_ID	The table ID of the active temporary table.
NAME	The name of the active temporary table.
N_COLS	The number of columns in the temporary table. The number always includes three hidden columns created by InnoDB (DB_ROW_ID, DB_TRX_ID, and DB_ROLL_PTR).
SPACE	The tablespace identifier (a numerical value) for the tablespace in which the temporary table resides. All InnoDB temporary tables reside in a shared temporary tablespace, as defined by innodb_temp_data_file_path. By default the shared temporary tablespace is named ibtmp1 and located in the data directory. Compressed temporary tables reside in separate per-table tablespaces located in the temporary file directory, as defined by tmpdir. The SPACE ID is always a non-zero value and is dynamically generated on server restart.

Example:

```
mysql> CREATE TEMPORARY TABLE t1 (c1 INT PRIMARY KEY) ENGINE=INNODB;
Query OK, 0 rows affected (0.01 sec)

mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_TEMP_TABLE_INFO\G
*************************** 1. row ***************************
TABLE_ID: 97
    NAME: #sql8c88_43_0
  N_COLS: 4
   SPACE: 76
1 row in set (0.00 sec)
```

Notes:

- This table is primarily useful for expert level monitoring.

- Use DESCRIBE or SHOW COLUMNS to view additional information about the columns of this table including data types and default values.
- You must have the PROCESS privilege to query this table.

24.31.28 The INFORMATION_SCHEMA INNODB_TRX Table

The INNODB_TRX table contains information about every transaction (excluding read-only transactions) currently executing inside InnoDB, including whether the transaction is waiting for a lock, when the transaction started, and the SQL statement the transaction is executing, if any.

Table 24.30 INNODB_TRX Columns

Column name	Description
TRX_ID	Unique transaction ID number, internal to InnoDB. These IDs are not created for transactions that are read only and nonlocking. See Section 8.5.3, “Optimizing InnoDB Read-Only Transactions” for details.
TRX_WEIGHT	The weight of a transaction, reflecting (but not necessarily the exact count of) the number of rows altered and the number of rows locked by the transaction. To resolve a deadlock, InnoDB selects the transaction with the smallest weight as the “victim” to roll back. Transactions that have changed non-transactional tables are considered heavier than others, regardless of the number of altered and locked rows.
TRX_STATE	Transaction execution state. Permitted values are RUNNING, LOCK WAIT, ROLLING BACK, and COMMITTING.
TRX_STARTED	Transaction start time.
TRX_REQUESTED_LOCK_ID	ID of the lock the transaction is currently waiting for, if TRX_STATE is LOCK WAIT; otherwise NULL. To obtain details about the lock, join this column with the ENGINE_LOCK_ID column of the Performance Schema data_locks table.
TRX_WAIT_STARTED	Time when the transaction started waiting on the lock, if TRX_STATE is LOCK WAIT; otherwise NULL.
TRX_MYSQL_THREAD_ID	MySQL thread ID. To obtain details about the thread, join this column with the ID column of the PROCESSLIST table, but see Section 15.14.2.3, “Persistence and Consistency of InnoDB Transaction and Locking Information”.
TRX_QUERY	The SQL statement that is being executed by the transaction.
TRX_OPERATION_STATE	The transaction's current operation, if any; otherwise NULL.
TRX_TABLES_IN_USE	The number of InnoDB tables used while processing the current SQL statement of this transaction.
TRX_TABLES_LOCKED	Number of InnoDB tables that the current SQL statement has row locks on. (Because these are row locks, not table locks, the tables can usually still be read from and written to by multiple transactions, despite some rows being locked.)
TRX_LOCK_STRUCTS	The number of locks reserved by the transaction.
TRX_LOCK_MEMORY_BYTES	Total size taken up by the lock structures of this transaction in memory.
TRX_ROWS_LOCKED	Approximate number or rows locked by this transaction. The value might include delete-marked rows that are physically present but not visible to the transaction.
TRX_ROWS_MODIFIED	The number of modified and inserted rows in this transaction.

Column name	Description
TRX_CONCURRENCY_TICKETS	A value indicating how much work the current transaction can do before being swapped out, as specified by the innodb_concurrency_tickets system variable.
TRX_ISOLATION_LEVEL	The isolation level of the current transaction.
TRX_UNIQUE_CHECKS	Whether unique checks are turned on or off for the current transaction. For example, they might be turned off during a bulk data load.
TRX_FOREIGN_KEY_CHECKS	Whether foreign key checks are turned on or off for the current transaction. For example, they might be turned off during a bulk data load.
TRX_LAST_FOREIGN_KEY_ERROR	Detailed error message for the last foreign key error, if any; otherwise NULL.
TRX_ADAPTIVE_HASH_LATCHED	Whether the adaptive hash index is locked by the current transaction. When the adaptive hash index search system is partitioned, a single transaction does not lock the entire adaptive hash index. Adaptive hash index partitioning is controlled by innodb_adaptive_hash_index_parts, which is set to 8 by default.
TRX_ADAPTIVE_HASH_TIMEOUT	Whether to relinquish the search latch immediately for the adaptive hash index, or reserve it across calls from MySQL. When there is no adaptive hash index contention, this value remains zero and statements reserve the latch until they finish. During times of contention, it counts down to zero, and statements release the latch immediately after each row lookup. When the adaptive hash index search system is partitioned (controlled by innodb_adaptive_hash_index_parts), the value remains 0.
TRX_IS_READ_ONLY	A value of 1 indicates the transaction is read only.
TRX_AUTOCOMMIT_NON_LOCKING	A value of 1 indicates the transaction is a SELECT statement that does not use the FOR UPDATE or LOCK IN SHARED MODE clauses, and is executing with autocommit enabled so that the transaction will only contain this one statement. When this column and TRX_IS_READ_ONLY are both 1, InnoDB optimizes the transaction to reduce the overhead associated with transactions that change table data.

Example:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.INNODB_TRX\G
*************************** 1. row ***************************
                    trx_id: 1510
                 trx_state: RUNNING
               trx_started: 2014-11-19 13:24:40
     trx_requested_lock_id: NULL
          trx_wait_started: NULL
                trx_weight: 586739
       trx_mysql_thread_id: 2
                 trx_query: DELETE FROM employees.salaries WHERE salary > 65000
       trx_operation_state: updating or deleting
         trx_tables_in_use: 1
         trx_tables_locked: 1
          trx_lock_structs: 3003
     trx_lock_memory_bytes: 450768
           trx_rows_locked: 1407513
         trx_rows_modified: 583736
   trx_concurrency_tickets: 0
       trx_isolation_level: REPEATABLE READ
         trx_unique_checks: 1
```

```
    trx_foreign_key_checks: 1
trx_last_foreign_key_error: NULL
 trx_adaptive_hash_latched: 0
 trx_adaptive_hash_timeout: 10000
          trx_is_read_only: 0
trx_autocommit_non_locking: 0
```

Notes:

- Use this table to help diagnose performance problems that occur during times of heavy concurrent load. Its contents are updated as described in Section 15.14.2.3, “Persistence and Consistency of InnoDB Transaction and Locking Information”.
- Use `DESCRIBE` or `SHOW COLUMNS` to view additional information about the columns of this table including data types and default values.
- You must have the `PROCESS` privilege to query this table.
- For usage information, see Section 15.14.2.1, “Using InnoDB Transaction and Locking Information”.

24.32 Connection-Control INFORMATION_SCHEMA Tables

The following sections describe the `INFORMATION_SCHEMA` tables associated with the `CONNECTION_CONTROL` plugin.

24.32.1 The INFORMATION_SCHEMA CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS Table

This table provides information about the current number of consecutive failed connection attempts per client user/host combination.

`CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` has these columns:

- `USERHOST`

  The user/host combination of a client that has failed connection attempts, in `'user_name'@'host_name'` format.

- `FAILED_ATTEMPTS`

  The current number of consecutive failed connection attempts for the `USERHOST` value. This counts all failed attempts, regardless of whether they were delayed. The number of attempts for which the server added a delay to its response is the difference between the `FAILED_ATTEMPTS` value and the `connection_control_failed_connections_threshold` system variable value.

Notes:

- The `CONNECTION_CONTROL_FAILED_LOGIN_ATTEMPTS` plugin must be activated for this table to be available, and the `CONNECTION_CONTROL` plugin must be activated or the table contents will always be empty. See Section 6.5.2, “The Connection-Control Plugins”.
- The table contains rows only for clients that have had one or more consecutive failed connection attempts without a subsequent successful attempt. When a client connects successfully, its failed-connection count is reset to zero and the server removes any row corresponding to the client.
- Assigning a value to the `connection_control_failed_connections_threshold` system variable at runtime resets all accumulated failed-connection counters to zero, which causes the table to become empty.

24.33 Extensions to SHOW Statements

Some extensions to SHOW statements accompany the implementation of INFORMATION_SCHEMA:

- SHOW can be used to get information about the structure of INFORMATION_SCHEMA itself.
- Several SHOW statements accept a WHERE clause that provides more flexibility in specifying which rows to display.

INFORMATION_SCHEMA is an information database, so its name is included in the output from SHOW DATABASES. Similarly, SHOW TABLES can be used with INFORMATION_SCHEMA to obtain a list of its tables:

```
mysql> SHOW TABLES FROM INFORMATION_SCHEMA;
+---------------------------------------+
| Tables_in_INFORMATION_SCHEMA          |
+---------------------------------------+
| CHARACTER_SETS                        |
| COLLATIONS                            |
| COLLATION_CHARACTER_SET_APPLICABILITY |
| COLUMNS                               |
| COLUMN_PRIVILEGES                     |
| ENGINES                               |
| EVENTS                                |
| FILES                                 |
| KEY_COLUMN_USAGE                      |
| PARTITIONS                            |
| PLUGINS                               |
| PROCESSLIST                           |
| REFERENTIAL_CONSTRAINTS               |
| ROUTINES                              |
| SCHEMATA                              |
| SCHEMA_PRIVILEGES                     |
| STATISTICS                            |
| TABLES                                |
| TABLE_CONSTRAINTS                     |
| TABLE_PRIVILEGES                      |
| TRIGGERS                              |
| USER_PRIVILEGES                       |
| VIEWS                                 |
+---------------------------------------+
```

SHOW COLUMNS and DESCRIBE can display information about the columns in individual INFORMATION_SCHEMA tables.

SHOW statements that accept a LIKE clause to limit the rows displayed also permit a WHERE clause that specifies more general conditions that selected rows must satisfy:

```
SHOW CHARACTER SET
SHOW COLLATION
SHOW COLUMNS
SHOW DATABASES
SHOW FUNCTION STATUS
SHOW INDEX
SHOW OPEN TABLES
SHOW PROCEDURE STATUS
SHOW STATUS
SHOW TABLE STATUS
SHOW TABLES
SHOW TRIGGERS
SHOW VARIABLES
```

The WHERE clause, if present, is evaluated against the column names displayed by the SHOW statement. For example, the SHOW CHARACTER SET statement produces these output columns:

```
mysql> SHOW CHARACTER SET;
+----------+-----------------------------+--------------------+--------+
```

```
| Charset  | Description                 | Default collation   | Maxlen |
+----------+-----------------------------+---------------------+--------+
| big5     | Big5 Traditional Chinese    | big5_chinese_ci     |      2 |
| dec8     | DEC West European           | dec8_swedish_ci     |      1 |
| cp850    | DOS West European           | cp850_general_ci    |      1 |
| hp8      | HP West European            | hp8_english_ci      |      1 |
| koi8r    | KOI8-R Relcom Russian       | koi8r_general_ci    |      1 |
| latin1   | cp1252 West European        | latin1_swedish_ci   |      1 |
| latin2   | ISO 8859-2 Central European | latin2_general_ci   |      1 |
...
```

To use a `WHERE` clause with `SHOW CHARACTER SET`, you would refer to those column names. As an example, the following statement displays information about character sets for which the default collation contains the string `'japanese'`:

```
mysql> SHOW CHARACTER SET WHERE `Default collation` LIKE '%japanese%';
+---------+---------------------------+---------------------+--------+
| Charset | Description               | Default collation   | Maxlen |
+---------+---------------------------+---------------------+--------+
| ujis    | EUC-JP Japanese           | ujis_japanese_ci    |      3 |
| sjis    | Shift-JIS Japanese        | sjis_japanese_ci    |      2 |
| cp932   | SJIS for Windows Japanese | cp932_japanese_ci   |      2 |
| eucjpms | UJIS for Windows Japanese | eucjpms_japanese_ci |      3 |
+---------+---------------------------+---------------------+--------+
```

This statement displays the multibyte character sets:

```
mysql> SHOW CHARACTER SET WHERE Maxlen > 1;
+---------+---------------------------+---------------------+--------+
| Charset | Description               | Default collation   | Maxlen |
+---------+---------------------------+---------------------+--------+
| big5    | Big5 Traditional Chinese  | big5_chinese_ci     |      2 |
| ujis    | EUC-JP Japanese           | ujis_japanese_ci    |      3 |
| sjis    | Shift-JIS Japanese        | sjis_japanese_ci    |      2 |
| euckr   | EUC-KR Korean             | euckr_korean_ci     |      2 |
| gb2312  | GB2312 Simplified Chinese | gb2312_chinese_ci   |      2 |
| gbk     | GBK Simplified Chinese    | gbk_chinese_ci      |      2 |
| utf8    | UTF-8 Unicode             | utf8_general_ci     |      3 |
| ucs2    | UCS-2 Unicode             | ucs2_general_ci     |      2 |
| cp932   | SJIS for Windows Japanese | cp932_japanese_ci   |      2 |
| eucjpms | UJIS for Windows Japanese | eucjpms_japanese_ci |      3 |
+---------+---------------------------+---------------------+--------+
```

Chapter 25 MySQL Performance Schema

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The MySQL Performance Schema is a feature for monitoring MySQL Server execution at a low level. The Performance Schema has these characteristics:

- The Performance Schema provides a way to inspect internal execution of the server at runtime. It is implemented using the `PERFORMANCE_SCHEMA` storage engine and the `performance_schema`

database. The Performance Schema focuses primarily on performance data. This differs from `INFORMATION_SCHEMA`, which serves for inspection of metadata.

- The Performance Schema monitors server events. An “event” is anything the server does that takes time and has been instrumented so that timing information can be collected. In general, an event could be a function call, a wait for the operating system, a stage of an SQL statement execution such as parsing or sorting, or an entire statement or group of statements. Event collection provides access to information about synchronization calls (such as for mutexes) file and table I/O, table locks, and so forth for the server and for several storage engines.
- Performance Schema events are distinct from events written to the server's binary log (which describe data modifications) and Event Scheduler events (which are a type of stored program).
- Performance Schema events are specific to a given instance of the MySQL Server. Performance Schema tables are considered local to the server, and changes to them are not replicated or written to the binary log.
- Current events are available, as well as event histories and summaries. This enables you to determine how many times instrumented activities were performed and how much time they took. Event information is available to show the activities of specific threads, or activity associated with particular objects such as a mutex or file.
- The `PERFORMANCE_SCHEMA` storage engine collects event data using “instrumentation points” in server source code.
- Collected events are stored in tables in the `performance_schema` database. These tables can be queried using `SELECT` statements like other tables.
- Performance Schema configuration can be modified dynamically by updating tables in the `performance_schema` database through SQL statements. Configuration changes affect data collection immediately.
- Tables in the Performance Schema are in-memory tables that use no persistent on-disk storage. The contents are repopulated beginning at server startup and discarded at server shutdown.
- Monitoring is available on all platforms supported by MySQL.

  Some limitations might apply: The types of timers might vary per platform. Instruments that apply to storage engines might not be implemented for all storage engines. Instrumentation of each third-party engine is the responsibility of the engine maintainer. See also Section C.8, “Restrictions on Performance Schema”.
- Data collection is implemented by modifying the server source code to add instrumentation. There are no separate threads associated with the Performance Schema, unlike other features such as replication or the Event Scheduler.

The Performance Schema is intended to provide access to useful information about server execution while having minimal impact on server performance. The implementation follows these design goals:

- Activating the Performance Schema causes no changes in server behavior. For example, it does not cause thread scheduling to change, and it does not cause query execution plans (as shown by `EXPLAIN`) to change.
- Server monitoring occurs continuously and unobtrusively with very little overhead. Activating the Performance Schema does not make the server unusable.
- The parser is unchanged. There are no new keywords or statements.
- Execution of server code proceeds normally even if the Performance Schema fails internally.
- When there is a choice between performing processing during event collection initially or during event retrieval later, priority is given to making collection faster. This is because collection is ongoing whereas retrieval is on demand and might never happen at all.

- Most Performance Schema tables have indexes, which gives the optimizer access to execution plans other than full table scans. For more information, see Section 8.2.4, “Optimizing Performance Schema Queries”.
- It is easy to add new instrumentation points.
- Instrumentation is versioned. If the instrumentation implementation changes, previously instrumented code will continue to work. This benefits developers of third-party plugins because it is not necessary to upgrade each plugin to stay synchronized with the latest Performance Schema changes.

Note

The MySQL `sys` schema is a set of objects that provides convenient access to data collected by the Performance Schema. The `sys` schema is installed by default. For usage instructions, see Chapter 26, *MySQL sys Schema*.

25.1 Performance Schema Quick Start

This section briefly introduces the Performance Schema with examples that show how to use it. For additional examples, see Section 25.18, “Using the Performance Schema to Diagnose Problems”.

The Performance Schema is enabled by default. To enable or disable it explicitly, start the server with the `performance_schema` variable set to an appropriate value. For example, use these lines in your `my.cnf` file:

```
[mysqld]
performance_schema=ON
```

When the server starts, it sees `performance_schema` and attempts to initialize the Performance Schema. To verify successful initialization, use this statement:

```
mysql> SHOW VARIABLES LIKE 'performance_schema';
+--------------------+-------+
| Variable_name      | Value |
+--------------------+-------+
| performance_schema | ON    |
+--------------------+-------+
```

A value of `ON` means that the Performance Schema initialized successfully and is ready for use. A value of `OFF` means that some error occurred. Check the server error log for information about what went wrong.

The Performance Schema is implemented as a storage engine, so you will see it listed in the output from the `INFORMATION_SCHEMA.ENGINES` table or the `SHOW ENGINES` statement:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.ENGINES
       WHERE ENGINE='PERFORMANCE_SCHEMA'\G
*************************** 1. row ***************************
      ENGINE: PERFORMANCE_SCHEMA
     SUPPORT: YES
     COMMENT: Performance Schema
TRANSACTIONS: NO
          XA: NO
  SAVEPOINTS: NO

mysql> SHOW ENGINES\G
...
      Engine: PERFORMANCE_SCHEMA
     Support: YES
     Comment: Performance Schema
Transactions: NO
          XA: NO
  Savepoints: NO
```

```
...
```

The PERFORMANCE_SCHEMA storage engine operates on tables in the performance_schema database. You can make performance_schema the default database so that references to its tables need not be qualified with the database name:

```
mysql> USE performance_schema;
```

Many examples in this chapter assume performance_schema as the default database.

Performance Schema tables are stored in the performance_schema database. Information about the structure of this database and its tables can be obtained, as for any other database, by selecting from the INFORMATION_SCHEMA database or by using SHOW statements. For example, use either of these statements to see what Performance Schema tables exist:

```
mysql> SELECT TABLE_NAME FROM INFORMATION_SCHEMA.TABLES
       WHERE TABLE_SCHEMA = 'performance_schema';
+----------------------------------------------------+
| TABLE_NAME                                         |
+----------------------------------------------------+
| accounts                                           |
| cond_instances                                     |
...
| events_stages_current                              |
| events_stages_history                              |
| events_stages_history_long                         |
| events_stages_summary_by_account_by_event_name     |
| events_stages_summary_by_host_by_event_name        |
| events_stages_summary_by_thread_by_event_name      |
| events_stages_summary_by_user_by_event_name        |
| events_stages_summary_global_by_event_name         |
| events_statements_current                          |
| events_statements_history                          |
| events_statements_history_long                     |
...
| file_instances                                     |
| file_summary_by_event_name                         |
| file_summary_by_instance                           |
| host_cache                                         |
| hosts                                              |
| memory_summary_by_account_by_event_name            |
| memory_summary_by_host_by_event_name               |
| memory_summary_by_thread_by_event_name             |
| memory_summary_by_user_by_event_name               |
| memory_summary_global_by_event_name                |
| metadata_locks                                     |
| mutex_instances                                    |
| objects_summary_global_by_type                     |
| performance_timers                                 |
| replication_connection_configuration               |
| replication_connection_status                      |
| replication_applier_configuration                  |
| replication_applier_status                         |
| replication_applier_status_by_coordinator          |
| replication_applier_status_by_worker               |
| rwlock_instances                                   |
| session_account_connect_attrs                      |
| session_connect_attrs                              |
| setup_actors                                       |
| setup_consumers                                    |
| setup_instruments                                  |
| setup_objects                                      |
| setup_timers                                       |
| socket_instances                                   |
| socket_summary_by_event_name                       |
| socket_summary_by_instance                         |
| table_handles                                      |
| table_io_waits_summary_by_index_usage              |
```

```
| table_io_waits_summary_by_table                     |
| table_lock_waits_summary_by_table                   |
| threads                                             |
| users                                               |
+-----------------------------------------------------+

mysql> SHOW TABLES FROM performance_schema;
+-----------------------------------------------------+
| Tables_in_performance_schema                        |
+-----------------------------------------------------+
| accounts                                            |
| cond_instances                                      |
| events_stages_current                               |
| events_stages_history                               |
| events_stages_history_long                          |
...
```

The number of Performance Schema tables increases over time as implementation of additional instrumentation proceeds.

The name of the `performance_schema` database is lowercase, as are the names of tables within it. Queries should specify the names in lowercase.

To see the structure of individual tables, use `SHOW CREATE TABLE`:

```
mysql> SHOW CREATE TABLE setup_timers\G
*************************** 1. row ***************************
       Table: setup_timers
Create Table: CREATE TABLE `setup_timers` (
  `NAME` varchar(64) NOT NULL,
  `TIMER_NAME` enum('CYCLE','NANOSECOND','MICROSECOND','MILLISECOND','TICK')
   NOT NULL
) ENGINE=PERFORMANCE_SCHEMA DEFAULT CHARSET=utf8
```

Table structure is also available by selecting from tables such as `INFORMATION_SCHEMA.COLUMNS` or by using statements such as `SHOW COLUMNS`.

Tables in the `performance_schema` database can be grouped according to the type of information in them: Current events, event histories and summaries, object instances, and setup (configuration) information. The following examples illustrate a few uses for these tables. For detailed information about the tables in each group, see Section 25.11, “Performance Schema Table Descriptions”.

Initially, not all instruments and consumers are enabled, so the performance schema does not collect all events. To turn all of these on and enable event timing, execute two statements (the row counts may differ depending on MySQL version):

```
mysql> UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES';
Query OK, 560 rows affected (0.04 sec)
mysql> UPDATE setup_consumers SET ENABLED = 'YES';
Query OK, 10 rows affected (0.00 sec)
```

To see what the server is doing at the moment, examine the `events_waits_current` table. It contains one row per thread showing each thread's most recent monitored event:

```
mysql> SELECT * FROM events_waits_current\G
*************************** 1. row ***************************
            THREAD_ID: 0
             EVENT_ID: 5523
         END_EVENT_ID: 5523
           EVENT_NAME: wait/synch/mutex/mysys/THR_LOCK::mutex
               SOURCE: thr_lock.c:525
          TIMER_START: 201660494489586
            TIMER_END: 201660494576112
           TIMER_WAIT: 86526
                SPINS: NULL
```

```
        OBJECT_SCHEMA: NULL
          OBJECT_NAME: NULL
           INDEX_NAME: NULL
          OBJECT_TYPE: NULL
OBJECT_INSTANCE_BEGIN: 142270668
     NESTING_EVENT_ID: NULL
   NESTING_EVENT_TYPE: NULL
            OPERATION: lock
      NUMBER_OF_BYTES: NULL
                FLAGS: 0
...
```

This event indicates that thread 0 was waiting for 86,526 picoseconds to acquire a lock on `THR_LOCK::mutex`, a mutex in the `mysys` subsystem. The first few columns provide the following information:

- The ID columns indicate which thread the event comes from and the event number.
- `EVENT_NAME` indicates what was instrumented and `SOURCE` indicates which source file contains the instrumented code.
- The timer columns show when the event started and stopped and how long it took. If an event is still in progress, the `TIMER_END` and `TIMER_WAIT` values are `NULL`. Timer values are approximate and expressed in picoseconds. For information about timers and event time collection, see Section 25.4.1, “Performance Schema Event Timing”.

The history tables contain the same kind of rows as the current-events table but have more rows and show what the server has been doing “recently” rather than “currently.” The `events_waits_history` and `events_waits_history_long` tables contain the most recent 10 events per thread and most recent 10,000 events, respectively. For example, to see information for recent events produced by thread 13, do this:

```
mysql> SELECT EVENT_ID, EVENT_NAME, TIMER_WAIT
       FROM events_waits_history WHERE THREAD_ID = 13
       ORDER BY EVENT_ID;
+----------+----------------------------------------+------------+
| EVENT_ID | EVENT_NAME                             | TIMER_WAIT |
+----------+----------------------------------------+------------+
|       86 | wait/synch/mutex/mysys/THR_LOCK::mutex |     686322 |
|       87 | wait/synch/mutex/mysys/THR_LOCK_malloc |     320535 |
|       88 | wait/synch/mutex/mysys/THR_LOCK_malloc |     339390 |
|       89 | wait/synch/mutex/mysys/THR_LOCK_malloc |     377100 |
|       90 | wait/synch/mutex/sql/LOCK_plugin       |     614673 |
|       91 | wait/synch/mutex/sql/LOCK_open         |     659925 |
|       92 | wait/synch/mutex/sql/THD::LOCK_thd_data |     494001 |
|       93 | wait/synch/mutex/mysys/THR_LOCK_malloc |     222489 |
|       94 | wait/synch/mutex/mysys/THR_LOCK_malloc |     214947 |
|       95 | wait/synch/mutex/mysys/LOCK_alarm      |     312993 |
+----------+----------------------------------------+------------+
```

As new events are added to a history table, older events are discarded if the table is full.

Summary tables provide aggregated information for all events over time. The tables in this group summarize event data in different ways. To see which instruments have been executed the most times or have taken the most wait time, sort the `events_waits_summary_global_by_event_name` table on the `COUNT_STAR` or `SUM_TIMER_WAIT` column, which correspond to a `COUNT(*)` or `SUM(TIMER_WAIT)` value, respectively, calculated over all events:

```
mysql> SELECT EVENT_NAME, COUNT_STAR
       FROM events_waits_summary_global_by_event_name
       ORDER BY COUNT_STAR DESC LIMIT 10;
+--------------------------------------------------+------------+
| EVENT_NAME                                       | COUNT_STAR |
+--------------------------------------------------+------------+
| wait/synch/mutex/mysys/THR_LOCK_malloc           |       6419 |
```

```
| wait/io/file/sql/FRM                                |        452 |
| wait/synch/mutex/sql/LOCK_plugin                    |        337 |
| wait/synch/mutex/mysys/THR_LOCK_open                |        187 |
| wait/synch/mutex/mysys/LOCK_alarm                   |        147 |
| wait/synch/mutex/sql/THD::LOCK_thd_data             |        115 |
| wait/io/file/myisam/kfile                           |        102 |
| wait/synch/mutex/sql/LOCK_global_system_variables |         89 |
| wait/synch/mutex/mysys/THR_LOCK::mutex              |         89 |
| wait/synch/mutex/sql/LOCK_open                      |         88 |
+-----------------------------------------------------+------------+

mysql> SELECT EVENT_NAME, SUM_TIMER_WAIT
       FROM events_waits_summary_global_by_event_name
       ORDER BY SUM_TIMER_WAIT DESC LIMIT 10;
+---------------------------------------+----------------+
| EVENT_NAME                            | SUM_TIMER_WAIT |
+---------------------------------------+----------------+
| wait/io/file/sql/MYSQL_LOG            |     1599816582 |
| wait/synch/mutex/mysys/THR_LOCK_malloc |     1530083250 |
| wait/io/file/sql/binlog_index         |     1385291934 |
| wait/io/file/sql/FRM                  |     1292823243 |
| wait/io/file/myisam/kfile             |      411193611 |
| wait/io/file/myisam/dfile             |      322401645 |
| wait/synch/mutex/mysys/LOCK_alarm     |      145126935 |
| wait/io/file/sql/casetest             |      104324715 |
| wait/synch/mutex/sql/LOCK_plugin      |       86027823 |
| wait/io/file/sql/pid                  |       72591750 |
+---------------------------------------+----------------+
```

These results show that the `THR_LOCK_malloc` mutex is “hot,” both in terms of how often it is used and amount of time that threads wait attempting to acquire it.

Note

The `THR_LOCK_malloc` mutex is used only in debug builds. In production builds it is not hot because it is nonexistent.

Instance tables document what types of objects are instrumented. An instrumented object, when used by the server, produces an event. These tables provide event names and explanatory notes or status information. For example, the `file_instances` table lists instances of instruments for file I/O operations and their associated files:

```
mysql> SELECT * FROM file_instances\G
*************************** 1. row ***************************
 FILE_NAME: /opt/mysql-log/60500/binlog.000007
EVENT_NAME: wait/io/file/sql/binlog
OPEN_COUNT: 0
*************************** 2. row ***************************
 FILE_NAME: /opt/mysql/60500/data/mysql/tables_priv.MYI
EVENT_NAME: wait/io/file/myisam/kfile
OPEN_COUNT: 1
*************************** 3. row ***************************
 FILE_NAME: /opt/mysql/60500/data/mysql/columns_priv.MYI
EVENT_NAME: wait/io/file/myisam/kfile
OPEN_COUNT: 1
...
```

Setup tables are used to configure and display monitoring characteristics. For example, to see which event timers are selected, query the `setup_timers` tables:

```
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | CYCLE       |
| stage       | NANOSECOND  |
| statement   | NANOSECOND  |
```

```
| transaction | NANOSECOND  |
+-------------+-------------+
```

`setup_instruments` lists the set of instruments for which events can be collected and shows which of them are enabled:

```
mysql> SELECT * FROM setup_instruments;
+-------------------------------------------------------------+---------+-------+
| NAME                                                        | ENABLED | TIMED |
+-------------------------------------------------------------+---------+-------+
...
| wait/synch/mutex/sql/LOCK_global_read_lock                  | YES     | YES   |
| wait/synch/mutex/sql/LOCK_global_system_variables           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_lock_db                           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_manager                           | YES     | YES   |
...
| wait/synch/rwlock/sql/LOCK_grant                            | YES     | YES   |
| wait/synch/rwlock/sql/LOGGER::LOCK_logger                   | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_connect                 | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_slave                   | YES     | YES   |
...
| wait/io/file/sql/binlog                                     | YES     | YES   |
| wait/io/file/sql/binlog_index                               | YES     | YES   |
| wait/io/file/sql/casetest                                   | YES     | YES   |
| wait/io/file/sql/dbopt                                      | YES     | YES   |
...
```

To understand how to interpret instrument names, see Section 25.6, “Performance Schema Instrument Naming Conventions”.

To control whether events are collected for an instrument, set its `ENABLED` value to `YES` or `NO`. For example:

```
mysql> UPDATE setup_instruments SET ENABLED = 'NO'
       WHERE NAME = 'wait/synch/mutex/sql/LOCK_mysql_create_db';
```

The Performance Schema uses collected events to update tables in the `performance_schema` database, which act as “consumers” of event information. The `setup_consumers` table lists the available consumers and which are enabled:

```
mysql> SELECT * FROM setup_consumers;
+--------------------------------+---------+
| NAME                           | ENABLED |
+--------------------------------+---------+
| events_stages_current          | NO      |
| events_stages_history          | NO      |
| events_stages_history_long     | NO      |
| events_statements_current      | YES     |
| events_statements_history      | YES     |
| events_statements_history_long | NO      |
| events_transactions_current    | YES     |
| events_transactions_history    | YES     |
| events_transactions_history_long | NO    |
| events_waits_current           | NO      |
| events_waits_history           | NO      |
| events_waits_history_long      | NO      |
| global_instrumentation         | YES     |
| thread_instrumentation         | YES     |
| statements_digest              | YES     |
+--------------------------------+---------+
```

To control whether the Performance Schema maintains a consumer as a destination for event information, set its `ENABLED` value.

For more information about the setup tables and how to use them to control event collection, see Section 25.4.2, “Performance Schema Event Filtering”.

There are some miscellaneous tables that do not fall into any of the previous groups. For example, `performance_timers` lists the available event timers and their characteristics. For information about timers, see Section 25.4.1, “Performance Schema Event Timing”.

25.2 Performance Schema Build Configuration

The Performance Schema is mandatory and always compiled in. It is possible to exclude certain parts of the Performance Schema instrumentation. For example, to exclude stage and statement instrumentation, do this:

```
shell> cmake . \
        -DDISABLE_PSI_STAGE=1 \
        -DDISABLE_PSI_STATEMENT=1
```

For more information, see the descriptions of the `DISABLE_PSI_`*`XXX`* `CMake` options in Section 2.8.4, “MySQL Source-Configuration Options”.

If you install MySQL over a previous installation that was configured without the Performance Schema (or with an older version of the Performance Schema that may not have all the current tables), run `mysql_upgrade` after starting the server to ensure that the `performance_schema` database exists with all current tables. Then restart the server. One indication that you need to do this is the presence of messages such as the following in the error log:

```
[ERROR] Native table 'performance_schema'.'events_waits_history'
has the wrong structure
[ERROR] Native table 'performance_schema'.'events_waits_history_long'
has the wrong structure
...
```

Because the Performance Schema is configured into the server at build time, a row for `PERFORMANCE_SCHEMA` appears in the output from `SHOW ENGINES`. This means that the Performance Schema is available, not that it is enabled. To enable it, you must do so at server startup, as described in the next section.

25.3 Performance Schema Startup Configuration

To use the MySQL Performance Schema, it must be enabled at server startup to enable event collection to occur.

The Performance Schema is enabled by default. To enable or disable it explicitly, start the server with the `performance_schema` variable set to an appropriate value. For example, use these lines in your `my.cnf` file:

```
[mysqld]
performance_schema=ON
```

If the server is unable to allocate any internal buffer during Performance Schema initialization, the Performance Schema disables itself and sets `performance_schema` to `OFF`, and the server runs without instrumentation.

The Performance Schema also permits instrument and consumer configuration at server startup.

To control an instrument at server startup, use an option of this form:

```
--performance-schema-instrument='instrument_name=value'
```

Here, *`instrument_name`* is an instrument name such as `wait/synch/mutex/sql/LOCK_open`, and *`value`* is one of these values:

- `OFF`, `FALSE`, or `0`: Disable the instrument
- `ON`, `TRUE`, or `1`: Enable and time the instrument
- `COUNTED`: Enable and count (rather than time) the instrument

Each `--performance-schema-instrument` option can specify only one instrument name, but multiple instances of the option can be given to configure multiple instruments. In addition, patterns are permitted in instrument names to configure instruments that match the pattern. To configure all condition synchronization instruments as enabled and counted, use this option:

```
--performance-schema-instrument='wait/synch/cond/%=COUNTED'
```

To disable all instruments, use this option:

```
--performance-schema-instrument='%=OFF'
```

Exception: The `memory/performance_schema/%` instruments are built in and cannot be disabled at startup.

Longer instrument name strings take precedence over shorter pattern names, regardless of order. For information about specifying patterns to select instruments, see Section 25.4.9, “Naming Instruments or Consumers for Filtering Operations”.

An unrecognized instrument name is ignored. It is possible that a plugin installed later may create the instrument, at which time the name is recognized and configured.

To control a consumer at server startup, use an option of this form:

```
--performance-schema-consumer-consumer_name=value
```

Here, *`consumer_name`* is a consumer name such as `events_waits_history`, and *`value`* is one of these values:

- `OFF`, `FALSE`, or `0`: Do not collect events for the consumer
- `ON`, `TRUE`, or `1`: Collect events for the consumer

For example, to enable the `events_waits_history` consumer, use this option:

```
--performance-schema-consumer-events-waits-history=ON
```

The permitted consumer names can be found by examining the `setup_consumers` table. Patterns are not permitted. Consumer names in the `setup_consumers` table use underscores, but for consumers set at startup, dashes and underscores within the name are equivalent.

The Performance Schema includes several system variables that provide configuration information:

```
mysql> SHOW VARIABLES LIKE 'perf%';
+----------------------------------------------------+---------+
| Variable_name                                      | Value   |
+----------------------------------------------------+---------+
| performance_schema                                 | ON      |
| performance_schema_accounts_size                   | 100     |
| performance_schema_digests_size                    | 200     |
| performance_schema_events_stages_history_long_size | 10000   |
| performance_schema_events_stages_history_size      | 10      |
| performance_schema_events_statements_history_long_size | 10000   |
| performance_schema_events_statements_history_size  | 10      |
| performance_schema_events_waits_history_long_size  | 10000   |
```

```
| performance_schema_events_waits_history_size           | 10      |
| performance_schema_hosts_size                          | 100     |
| performance_schema_max_cond_classes                    | 80      |
| performance_schema_max_cond_instances                  | 1000    |
...
```

The `performance_schema` variable is `ON` or `OFF` to indicate whether the Performance Schema is enabled or disabled. The other variables indicate table sizes (number of rows) or memory allocation values.

Note

With the Performance Schema enabled, the number of Performance Schema instances affects the server memory footprint, perhaps to a large extent. The Performance Schema autoscales many parameters to use memory only as required; see Section 25.16, “The Performance Schema Memory-Allocation Model”.

To change the value of Performance Schema system variables, set them at server startup. For example, put the following lines in a `my.cnf` file to change the sizes of the history tables for wait events:

```
[mysqld]
performance_schema
performance_schema_events_waits_history_size=20
performance_schema_events_waits_history_long_size=15000
```

The Performance Schema automatically sizes the values of several of its parameters at server startup if they are not set explicitly. For example, it sizes the parameters that control the sizes of the events waits tables this way. The Performance Schema allocates memory incrementally, scaling its memory use to actual server load, instead of allocating all the memory it needs during server startup. Consequently, many sizing parameters need not be set at all. To see which parameters are autosized or autoscaled, use `mysqld --verbose --help` and examine the option descriptions, or see Section 25.14, “Performance Schema System Variables”.

For each autosized parameter that is not set at server startup (or is set to −1), the Performance Schema determines how to set its value based on the value of the following system values, which are considered as “hints” about how you have configured your MySQL server:

```
max_connections
open_files_limit
table_definition_cache
table_open_cache
```

To override autosizing or autoscaling for a given parameter, set it to a value other than −1 at startup. In this case, the Performance Schema assigns it the specified value.

At runtime, `SHOW VARIABLES` displays the actual values that autosized parameters were set to. Autoscaled parameters display with a value of −1.

If the Performance Schema is disabled, its autosized and autoscaled parameters remain set to −1 and `SHOW VARIABLES` displays −1.

25.4 Performance Schema Runtime Configuration

Specific Performance Schema features can be enabled at runtime to control which types of event collection occur.

Performance Schema setup tables contain information about monitoring configuration:

```
mysql> SELECT TABLE_NAME FROM INFORMATION_SCHEMA.TABLES
       WHERE TABLE_SCHEMA = 'performance_schema'
       AND TABLE_NAME LIKE 'setup%';
+-------------------+
| TABLE_NAME        |
+-------------------+
| setup_actors      |
| setup_consumers   |
| setup_instruments |
| setup_objects     |
| setup_timers      |
+-------------------+
```

You can examine the contents of these tables to obtain information about Performance Schema monitoring characteristics. If you have the `UPDATE` privilege, you can change Performance Schema operation by modifying setup tables to affect how monitoring occurs. For additional details about these tables, see Section 25.11.2, “Performance Schema Setup Tables”.

To see which event timers are selected, query the `setup_timers` tables:

```
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | CYCLE       |
| stage       | NANOSECOND  |
| statement   | NANOSECOND  |
| transaction | NANOSECOND  |
+-------------+-------------+
```

The `NAME` value indicates the type of instrument to which the timer applies, and `TIMER_NAME` indicates which timer applies to those instruments. The timer applies to instruments where their name begins with a component matching the `NAME` value.

To change the timer, update the `NAME` value. For example, to use the `NANOSECOND` timer for the `wait` timer:

```
mysql> UPDATE setup_timers SET TIMER_NAME = 'NANOSECOND'
       WHERE NAME = 'wait';
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | NANOSECOND  |
| stage       | NANOSECOND  |
| statement   | NANOSECOND  |
| transaction | NANOSECOND  |
+-------------+-------------+
```

For discussion of timers, see Section 25.4.1, “Performance Schema Event Timing”.

The `setup_instruments` and `setup_consumers` tables list the instruments for which events can be collected and the types of consumers for which event information actually is collected, respectively. Other setup tables enable further modification of the monitoring configuration. Section 25.4.2, “Performance Schema Event Filtering”, discusses how you can modify these tables to affect event collection.

If there are Performance Schema configuration changes that must be made at runtime using SQL statements and you would like these changes to take effect each time the server starts, put the statements in a file and start the server with the `--init-file=file_name` option. This strategy can also be useful if you have multiple monitoring configurations, each tailored to produce a different kind of monitoring, such as casual server health monitoring, incident investigation, application behavior

troubleshooting, and so forth. Put the statements for each monitoring configuration into their own file and specify the appropriate file as the `--init-file` argument when you start the server.

25.4.1 Performance Schema Event Timing

Events are collected by means of instrumentation added to the server source code. Instruments time events, which is how the Performance Schema provides an idea of how long events take. It is also possible to configure instruments not to collect timing information. This section discusses the available timers and their characteristics, and how timing values are represented in events.

Performance Schema Timers

Two Performance Schema tables provide timer information:

- `performance_timers` lists the available timers and their characteristics.
- `setup_timers` indicates which timers are used for which instruments.

Each timer row in `setup_timers` must refer to one of the timers listed in `performance_timers`.

Timers vary in precision and amount of overhead. To see what timers are available and their characteristics, check the `performance_timers` table:

```
mysql> SELECT * FROM performance_timers;
+-------------+-----------------+------------------+----------------+
| TIMER_NAME  | TIMER_FREQUENCY | TIMER_RESOLUTION | TIMER_OVERHEAD |
+-------------+-----------------+------------------+----------------+
| CYCLE       |      2389029850 |                1 |             72 |
| NANOSECOND  |      1000000000 |                1 |            112 |
| MICROSECOND |         1000000 |                1 |            136 |
| MILLISECOND |            1036 |                1 |            168 |
| TICK        |             105 |                1 |           2416 |
+-------------+-----------------+------------------+----------------+
```

The columns have these meanings:

- The `TIMER_NAME` column shows the names of the available timers. `CYCLE` refers to the timer that is based on the CPU (processor) cycle counter. The timers in `setup_timers` that you can use are those that do not have `NULL` in the other columns. If the values associated with a given timer name are `NULL`, that timer is not supported on your platform.
- `TIMER_FREQUENCY` indicates the number of timer units per second. For a cycle timer, the frequency is generally related to the CPU speed. The value shown was obtained on a system with a 2.4GHz processor. The other timers are based on fixed fractions of seconds. For `TICK`, the frequency may vary by platform (for example, some use 100 ticks/second, others 1000 ticks/second).
- `TIMER_RESOLUTION` indicates the number of timer units by which timer values increase at a time. If a timer has a resolution of 10, its value increases by 10 each time.
- `TIMER_OVERHEAD` is the minimal number of cycles of overhead to obtain one timing with the given timer. The overhead per event is twice the value displayed because the timer is invoked at the beginning and end of the event.

To see which timers are in effect or to change timers, access the `setup_timers` table:

```
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | CYCLE       |
| stage       | NANOSECOND  |
```

```
| statement   | NANOSECOND  |
| transaction | NANOSECOND  |
+-------------+-------------+

mysql> UPDATE setup_timers SET TIMER_NAME = 'MICROSECOND'
       WHERE NAME = 'idle';
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | CYCLE       |
| stage       | NANOSECOND  |
| statement   | NANOSECOND  |
| transaction | NANOSECOND  |
+-------------+-------------+
```

By default, the Performance Schema uses the best timer available for each instrument type, but you can select a different one.

To time wait events, the most important criterion is to reduce overhead, at the possible expense of the timer accuracy, so using the `CYCLE` timer is the best.

The time a statement (or stage) takes to execute is in general orders of magnitude larger than the time it takes to execute a single wait. To time statements, the most important criterion is to have an accurate measure, which is not affected by changes in processor frequency, so using a timer which is not based on cycles is the best. The default timer for statements is `NANOSECOND`. The extra “overhead” compared to the `CYCLE` timer is not significant, because the overhead caused by calling a timer twice (once when the statement starts, once when it ends) is orders of magnitude less compared to the CPU time used to execute the statement itself. Using the `CYCLE` timer has no benefit here, only drawbacks.

The precision offered by the cycle counter depends on processor speed. If the processor runs at 1 GHz (one billion cycles/second) or higher, the cycle counter delivers sub-nanosecond precision. Using the cycle counter is much cheaper than getting the actual time of day. For example, the standard `gettimeofday()` function can take hundreds of cycles, which is an unacceptable overhead for data gathering that may occur thousands or millions of times per second.

Cycle counters also have disadvantages:

- End users expect to see timings in wall-clock units, such as fractions of a second. Converting from cycles to fractions of seconds can be expensive. For this reason, the conversion is a quick and fairly rough multiplication operation.
- Processor cycle rate might change, such as when a laptop goes into power-saving mode or when a CPU slows down to reduce heat generation. If a processor's cycle rate fluctuates, conversion from cycles to real-time units is subject to error.
- Cycle counters might be unreliable or unavailable depending on the processor or the operating system. For example, on Pentiums, the instruction is `RDTSC` (an assembly-language rather than a C instruction) and it is theoretically possible for the operating system to prevent user-mode programs from using it.
- Some processor details related to out-of-order execution or multiprocessor synchronization might cause the counter to seem fast or slow by up to 1000 cycles.

MySQL works with cycle counters on x386 (Windows, OS X, Linux, Solaris, and other Unix flavors), PowerPC, and IA-64.

Performance Schema Timer Representation in Events

Rows in Performance Schema tables that store current events and historical events have three columns to represent timing information: `TIMER_START` and `TIMER_END` indicate when an event started and finished, and `TIMER_WAIT` indicates event duration.

The `setup_instruments` table has an `ENABLED` column to indicate the instruments for which to collect events. The table also has a `TIMED` column to indicate which instruments are timed. If an instrument is not enabled, it produces no events. If an enabled instrument is not timed, events produced by the instrument have `NULL` for the `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` timer values. This in turn causes those values to be ignored when calculating the sum, minimum, maximum, and average time values in summary tables.

Internally, times within events are stored in units given by the timer in effect when event timing begins. For display when events are retrieved from Performance Schema tables, times are shown in picoseconds (trillionths of a second) to normalize them to a standard unit, regardless of which timer is selected.

Modifications to the `setup_timers` table affect monitoring immediately. Events already in progress may use the original timer for the begin time and the new timer for the end time. To avoid unpredictable results after you make timer changes, use `TRUNCATE TABLE` to reset Performance Schema statistics.

The timer baseline (“time zero”) occurs at Performance Schema initialization during server startup. `TIMER_START` and `TIMER_END` values in events represent picoseconds since the baseline. `TIMER_WAIT` values are durations in picoseconds.

Picosecond values in events are approximate. Their accuracy is subject to the usual forms of error associated with conversion from one unit to another. If the `CYCLE` timer is used and the processor rate varies, there might be drift. For these reasons, it is not reasonable to look at the `TIMER_START` value for an event as an accurate measure of time elapsed since server startup. On the other hand, it is reasonable to use `TIMER_START` or `TIMER_WAIT` values in `ORDER BY` clauses to order events by start time or duration.

The choice of picoseconds in events rather than a value such as microseconds has a performance basis. One implementation goal was to show results in a uniform time unit, regardless of the timer. In an ideal world this time unit would look like a wall-clock unit and be reasonably precise; in other words, microseconds. But to convert cycles or nanoseconds to microseconds, it would be necessary to perform a division for every instrumentation. Division is expensive on many platforms. Multiplication is not expensive, so that is what is used. Therefore, the time unit is an integer multiple of the highest possible `TIMER_FREQUENCY` value, using a multiplier large enough to ensure that there is no major precision loss. The result is that the time unit is “picoseconds.” This precision is spurious, but the decision enables overhead to be minimized.

While a wait, stage, statement, or transaction event is executing, the respective current-event tables display current-event timing information:

```
events_waits_current
events_stages_current
events_statements_current
events_transactions_current
```

To make it possible to determine how how long a not-yet-completed event has been running, the timer columns are set as follows:

- `TIMER_START` is populated (unchanged from previous behavior)
- `TIMER_END` is populated with the current timer value
- `TIMER_WAIT` is populated with the time elapsed so far (`TIMER_END` − `TIMER_START`)

Events that have not yet completed have an `END_EVENT_ID` value of `NULL`. To assess time elapsed so far for an event, use the `TIMER_WAIT` column. Therefore, to identify events that have not yet completed and have taken longer than *`N`* picoseconds thus far, monitoring applications can use this expression in queries:

```
WHERE END_EVENT_ID IS NULL AND TIMER_WAIT > N
```

Event identification as just described assumes that the corresponding instruments have `ENABLED` and `TIMED` set to `YES` and that the relevant consumers are enabled.

25.4.2 Performance Schema Event Filtering

Events are processed in a producer/consumer fashion:

- Instrumented code is the source for events and produces events to be collected. The `setup_instruments` table lists the instruments for which events can be collected, whether they are enabled, and (for enabled instruments) whether to collect timing information:

```
mysql> SELECT * FROM setup_instruments;
+-------------------------------------------------------------+---------+-------+
| NAME                                                        | ENABLED | TIMED |
+-------------------------------------------------------------+---------+-------+
...
| wait/synch/mutex/sql/LOCK_global_read_lock                  | YES     | YES   |
| wait/synch/mutex/sql/LOCK_global_system_variables           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_lock_db                           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_manager                           | YES     | YES   |
...
```

  The `setup_instruments` table provides the most basic form of control over event production. To further refine event production based on the type of object or thread being monitored, other tables may be used as described in Section 25.4.3, “Event Pre-Filtering”.

- Performance Schema tables are the destinations for events and consume events. The `setup_consumers` table lists the types of consumers to which event information can be sent and whether they are enabled:

```
mysql> SELECT * FROM setup_consumers;
+---------------------------------+---------+
| NAME                            | ENABLED |
+---------------------------------+---------+
| events_stages_current           | NO      |
| events_stages_history           | NO      |
| events_stages_history_long      | NO      |
| events_statements_current       | YES     |
| events_statements_history       | YES     |
| events_statements_history_long  | NO      |
| events_transactions_current     | YES     |
| events_transactions_history     | YES     |
| events_transactions_history_long | NO      |
| events_waits_current            | NO      |
| events_waits_history            | NO      |
| events_waits_history_long       | NO      |
| global_instrumentation          | YES     |
| thread_instrumentation          | YES     |
| statements_digest               | YES     |
+---------------------------------+---------+
```

Filtering can be done at different stages of performance monitoring:

- **Pre-filtering.** This is done by modifying Performance Schema configuration so that only certain types of events are collected from producers, and collected events update only certain consumers. To do this, enable or disable instruments or consumers. Pre-filtering is done by the Performance Schema and has a global effect that applies to all users.

  Reasons to use pre-filtering:

  - To reduce overhead. Performance Schema overhead should be minimal even with all instruments enabled, but perhaps you want to reduce it further. Or you do not care about timing events and want to disable the timing code to eliminate timing overhead.

  - To avoid filling the current-events or history tables with events in which you have no interest. Pre-filtering leaves more “room” in these tables for instances of rows for enabled instrument types. If you enable only file instruments with pre-filtering, no rows are collected for nonfile instruments. With post-filtering, nonfile events are collected, leaving fewer rows for file events.

  - To avoid maintaining some kinds of event tables. If you disable a consumer, the server does not spend time maintaining destinations for that consumer. For example, if you do not care about event histories, you can disable the history table consumers to improve performance.

- **Post-filtering.** This involves the use of `WHERE` clauses in queries that select information from Performance Schema tables, to specify which of the available events you want to see. Post-filtering is performed on a per-user basis because individual users select which of the available events are of interest.

  Reasons to use post-filtering:

  - To avoid making decisions for individual users about which event information is of interest.

  - To use the Performance Schema to investigate a performance issue when the restrictions to impose using pre-filtering are not known in advance.

The following sections provide more detail about pre-filtering and provide guidelines for naming instruments or consumers in filtering operations. For information about writing queries to retrieve information (post-filtering), see Section 25.5, “Performance Schema Queries”.

25.4.3 Event Pre-Filtering

Pre-filtering is done by the Performance Schema and has a global effect that applies to all users. Pre-filtering can be applied to either the producer or consumer stage of event processing:

- To configure pre-filtering at the producer stage, several tables can be used:

  - `setup_instruments` indicates which instruments are available. An instrument disabled in this table produces no events regardless of the contents of the other production-related setup tables. An instrument enabled in this table is permitted to produce events, subject to the contents of the other tables.

  - `setup_objects` controls whether the Performance Schema monitors particular table and stored program objects.

  - `threads` indicates whether monitoring is enabled for each server thread.

  - `setup_actors` determines the initial monitoring state for new foreground threads.

- To configure pre-filtering at the consumer stage, modify the `setup_consumers` table. This determines the destinations to which events are sent. `setup_consumers` also implicitly affects event production. If a given event will not be sent to any destination (that is, will not be consumed), the Performance Schema does not produce it.

Modifications to any of these tables affect monitoring immediately, with the exception that modifications to the `setup_actors` table affect only foreground threads created subsequent to the modification, not existing threads.

When you change the monitoring configuration, the Performance Schema does not flush the history tables. Events already collected remain in the current-events and history tables until displaced by newer events. If you disable instruments, you might need to wait a while before events for them are displaced by newer events of interest. Alternatively, use `TRUNCATE TABLE` to empty the history tables.

After making instrumentation changes, you might want to truncate the summary tables. Generally, the effect is to reset the summary columns to 0 or `NULL`, not to remove rows. This enables you to clear

collected values and restart aggregation. That might be useful, for example, after you have made a runtime configuration change. Exceptions to this truncation behavior are noted in individual summary table sections.

The following sections describe how to use specific tables to control Performance Schema pre-filtering.

25.4.4 Pre-Filtering by Instrument

The `setup_instruments` table lists the available instruments:

```
mysql> SELECT * FROM setup_instruments;
+-------------------------------------------------------------+---------+-------+
| NAME                                                        | ENABLED | TIMED |
+-------------------------------------------------------------+---------+-------+
...
| wait/synch/mutex/sql/LOCK_global_read_lock                  | YES     | YES   |
| wait/synch/mutex/sql/LOCK_global_system_variables           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_lock_db                           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_manager                           | YES     | YES   |
...
| wait/synch/rwlock/sql/LOCK_grant                            | YES     | YES   |
| wait/synch/rwlock/sql/LOGGER::LOCK_logger                   | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_connect                 | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_slave                   | YES     | YES   |
...
| wait/io/file/sql/binlog                                     | YES     | YES   |
| wait/io/file/sql/binlog_index                               | YES     | YES   |
| wait/io/file/sql/casetest                                   | YES     | YES   |
| wait/io/file/sql/dbopt                                      | YES     | YES   |
...
```

To control whether an instrument is enabled, set its `ENABLED` column to `YES` or `NO`. To configure whether to collect timing information for an enabled instrument, set its `TIMED` value to `YES` or `NO`. Setting the `TIMED` column affects Performance Schema table contents as described in Section 25.4.1, “Performance Schema Event Timing”.

Modifications to most `setup_instruments` rows affect monitoring immediately. For some instruments, modifications are effective only at server startup; changing them at runtime has no effect. This affects primarily mutexes, conditions, and rwlocks in the server, although there may be other instruments for which this is true.

The `setup_instruments` table provides the most basic form of control over event production. To further refine event production based on the type of object or thread being monitored, other tables may be used as described in Section 25.4.3, “Event Pre-Filtering”.

The following examples demonstrate possible operations on the `setup_instruments` table. These changes, like other pre-filtering operations, affect all users. Some of these queries use the `LIKE` operator and a pattern match instrument names. For additional information about specifying patterns to select instruments, see Section 25.4.9, “Naming Instruments or Consumers for Filtering Operations”.

- Disable all instruments:

  ```
  mysql> UPDATE setup_instruments SET ENABLED = 'NO';
  ```

  Now no events will be collected.

- Disable all file instruments, adding them to the current set of disabled instruments:

  ```
  mysql> UPDATE setup_instruments SET ENABLED = 'NO'
         WHERE NAME LIKE 'wait/io/file/%';
  ```

- Disable only file instruments, enable all other instruments:

```
mysql> UPDATE setup_instruments
       SET ENABLED = IF(NAME LIKE 'wait/io/file/%', 'NO', 'YES');
```

- Enable all but those instruments in the `mysys` library:

```
mysql> UPDATE setup_instruments
       SET ENABLED = CASE WHEN NAME LIKE '%/mysys/%' THEN 'YES' ELSE 'NO' END;
```

- Disable a specific instrument:

```
mysql> UPDATE setup_instruments SET ENABLED = 'NO'
       WHERE NAME = 'wait/synch/mutex/mysys/TMPDIR_mutex';
```

- To toggle the state of an instrument, “flip” its `ENABLED` value:

```
mysql> UPDATE setup_instruments
       SET ENABLED = IF(ENABLED = 'YES', 'NO', 'YES')
       WHERE NAME = 'wait/synch/mutex/mysys/TMPDIR_mutex';
```

- Disable timing for all events:

```
mysql> UPDATE setup_instruments SET TIMED = 'NO';
```

25.4.5 Pre-Filtering by Object

The `setup_objects` table controls whether the Performance Schema monitors particular table and stored program objects. The initial `setup_objects` contents look like this:

```
mysql> SELECT * FROM setup_objects;
+-------------+--------------------+-------------+---------+-------+
| OBJECT_TYPE | OBJECT_SCHEMA      | OBJECT_NAME | ENABLED | TIMED |
+-------------+--------------------+-------------+---------+-------+
| EVENT       | mysql              | %           | NO      | NO    |
| EVENT       | performance_schema | %           | NO      | NO    |
| EVENT       | information_schema | %           | NO      | NO    |
| EVENT       | %                  | %           | YES     | YES   |
| FUNCTION    | mysql              | %           | NO      | NO    |
| FUNCTION    | performance_schema | %           | NO      | NO    |
| FUNCTION    | information_schema | %           | NO      | NO    |
| FUNCTION    | %                  | %           | YES     | YES   |
| PROCEDURE   | mysql              | %           | NO      | NO    |
| PROCEDURE   | performance_schema | %           | NO      | NO    |
| PROCEDURE   | information_schema | %           | NO      | NO    |
| PROCEDURE   | %                  | %           | YES     | YES   |
| TABLE       | mysql              | %           | NO      | NO    |
| TABLE       | performance_schema | %           | NO      | NO    |
| TABLE       | information_schema | %           | NO      | NO    |
| TABLE       | %                  | %           | YES     | YES   |
| TRIGGER     | mysql              | %           | NO      | NO    |
| TRIGGER     | performance_schema | %           | NO      | NO    |
| TRIGGER     | information_schema | %           | NO      | NO    |
| TRIGGER     | %                  | %           | YES     | YES   |
+-------------+--------------------+-------------+---------+-------+
```

Modifications to the `setup_objects` table affect object monitoring immediately.

The `OBJECT_TYPE` column indicates the type of object to which a row applies. `TABLE` filtering affects table I/O events (`wait/io/table/sql/handler` instrument) and table lock events (`wait/lock/table/sql/handler` instrument).

The `OBJECT_SCHEMA` and `OBJECT_NAME` columns should contain a literal schema or object name, or `'%'` to match any name.

The ENABLED column indicates whether matching objects are monitored, and TIMED indicates whether to collect timing information. Setting the TIMED column affects Performance Schema table contents as described in Section 25.4.1, “Performance Schema Event Timing”.

The effect of the default object configuration is to instrument all objects except those in the mysql, INFORMATION_SCHEMA, and performance_schema databases. (Tables in the INFORMATION_SCHEMA database are not instrumented regardless of the contents of setup_objects; the row for information_schema.% simply makes this default explicit.)

When the Performance Schema checks for a match in setup_objects, it tries to find more specific matches first. For rows that match a given OBJECT_TYPE, the Performance Schema checks rows in this order:

- Rows with OBJECT_SCHEMA='*literal*' and OBJECT_NAME='*literal*'.
- Rows with OBJECT_SCHEMA='*literal*' and OBJECT_NAME='%'.
- Rows with OBJECT_SCHEMA='%' and OBJECT_NAME='%'.

For example, with a table db1.t1, the Performance Schema looks in TABLE rows for a match for 'db1' and 't1', then for 'db1' and '%', then for '%' and '%'. The order in which matching occurs matters because different matching setup_objects rows can have different ENABLED and TIMED values.

For table-related events, the Performance Schema combines the contents of setup_objects with setup_instruments to determine whether to enable instruments and whether to time enabled instruments:

- For tables that match a row in setup_objects, table instruments produce events only if ENABLED is YES in both setup_instruments and setup_objects.
- The TIMED values in the two tables are combined, so that timing information is collected only when both values are YES.

For stored program objects, the Performance Schema takes the ENABLED and TIMED columns directly from the setup_objects row. There is no combining of values with setup_instruments.

Suppose that setup_objects contains the following TABLE rows that apply to db1, db2, and db3:

```
+-------------+---------------+-------------+---------+-------+
| OBJECT_TYPE | OBJECT_SCHEMA | OBJECT_NAME | ENABLED | TIMED |
+-------------+---------------+-------------+---------+-------+
| TABLE       | db1           | t1          | YES     | YES   |
| TABLE       | db1           | t2          | NO      | NO    |
| TABLE       | db2           | %           | YES     | YES   |
| TABLE       | db3           | %           | NO      | NO    |
| TABLE       | %             | %           | YES     | YES   |
+-------------+---------------+-------------+---------+-------+
```

If an object-related instrument in setup_instruments has an ENABLED value of NO, events for the object are not monitored. If the ENABLED value is YES, event monitoring occurs according to the ENABLED value in the relevant setup_objects row:

- db1.t1 events are monitored
- db1.t2 events are not monitored
- db2.t3 events are monitored
- db3.t4 events are not monitored
- db4.t5 events are monitored

Similar logic applies for combining the `TIMED` columns from the `setup_instruments` and `setup_objects` tables to determine whether to collect event timing information.

If a persistent table and a temporary table have the same name, matching against `setup_objects` rows occurs the same way for both. It is not possible to enable monitoring for one table but not the other. However, each table is instrumented separately.

25.4.6 Pre-Filtering by Thread

The `threads` table contains a row for each server thread. Each row contains information about a thread and indicates whether monitoring is enabled for it. For the Performance Schema to monitor a thread, these things must be true:

- The `thread_instrumentation` consumer in the `setup_consumers` table must be `YES`.
- The `threads.INSTRUMENTED` column must be `YES`.
- Monitoring occurs only for those thread events produced from instruments that are enabled in the `setup_instruments` table.

The `threads` table also indicates for each server thread whether to perform historical event logging. This includes wait, stage, statement, and transaction events and affects logging to these tables:

```
events_waits_history
events_waits_history_long
events_stages_history
events_stages_history_long
events_statements_history
events_statements_history_long
events_transactions_history
events_transactions_history_long
```

For historical event logging to occur, these things must be true:

- The appropriate history-related consumers in the `setup_consumers` table must be enabled. For example, wait event logging in the `events_waits_history` and `events_waits_history_long` tables requires the corresponding `events_waits_history` and `events_waits_history_long` consumers to be `YES`.
- The `threads.HISTORY` column must be `YES`.
- Logging occurs only for those thread events produced from instruments that are enabled in the `setup_instruments` table.

For foreground threads (resulting from client connections), the initial values of the `INSTRUMENTED` and `HISTORY` columns in `threads` table rows are determined by whether the user account associated with a thread matches any row in the `setup_actors` table. The values come from the `ENABLED` and `HISTORY` columns of the matching `setup_actors` table row.

For background threads, there is no associated user. `INSTRUMENTED` and `HISTORY` are `YES` by default and `setup_actors` is not consulted.

The initial `setup_actors` contents look like this:

```
mysql> SELECT * FROM setup_actors;
+------+------+------+---------+---------+
| HOST | USER | ROLE | ENABLED | HISTORY |
+------+------+------+---------+---------+
| %    | %    | %    | YES     | YES     |
+------+------+------+---------+---------+
```

The `HOST` and `USER` columns should contain a literal host or user name, or `'%'` to match any name.

The ENABLED and HISTORY columns indicate whether to enable instrumentation and historical event logging for matching threads, subject to the other conditions described previously.

When the Performance Schema checks for a match for each new foreground thread in setup_actors, it tries to find more specific matches first, using the USER and HOST columns (ROLE is unused):

- Rows with USER='*literal*' and HOST='*literal*'.
- Rows with USER='*literal*' and HOST='%'.
- Rows with USER='%' and HOST='*literal*'.
- Rows with USER='%' and HOST='%'.

The order in which matching occurs matters because different matching setup_actors rows can have different USER and HOST values. This enables instrumenting and historical event logging to be applied selectively per host, user, or account (user and host combination), based on the ENABLED and HISTORY column values:

- When the best match is a row with ENABLED=YES, the INSTRUMENTED value for the thread becomes YES. When the best match is a row with HISTORY=YES, the HISTORY value for the thread becomes YES.
- When the best match is a row with ENABLED=NO, the INSTRUMENTED value for the thread becomes NO. When the best match is a row with HISTORY=NO, the HISTORY value for the thread becomes NO.
- When no match is found, the INSTRUMENTED and HISTORY values for the thread become NO.

The ENABLED and HISTORY columns in setup_actors rows can be set to YES or NO independent of one another. This means you can enable instrumentation separately from whether you collect historical events.

By default, monitoring and historical event collection are enabled for all new foreground threads because the setup_actors table initially contains a row with '%' for both HOST and USER. To perform more limited matching such as to enable monitoring only for some foreground threads, you must change this row because it matches any connection, and add rows for more specific HOST/USER combinations.

Suppose that you modify setup_actors as follows:

```
UPDATE setup_actors SET ENABLED = 'NO', HISTORY = 'NO'
WHERE HOST = '%' AND USER = '%';
INSERT INTO setup_actors (HOST,USER,ROLE,ENABLED,HISTORY)
VALUES('localhost','joe','%','YES','YES');
INSERT INTO setup_actors (HOST,USER,ROLE,ENABLED,HISTORY)
VALUES('hosta.example.com','joe','%','YES','NO');
INSERT INTO setup_actors (HOST,USER,ROLE,ENABLED,HISTORY)
VALUES('%','sam','%','NO','YES');
```

The UPDATE statement changes the default match to disable instrumentation and historical event collection. The INSERT statements add rows for more specific matches.

Now the Performance Schema determines how to set the INSTRUMENTED and HISTORY values for new connection threads as follows:

- If joe connects from the local host, the connection matches the first inserted row. The INSTRUMENTED and HISTORY values for the thread become YES.
- If joe connects from hosta.example.com, the connection matches the second inserted row. The INSTRUMENTED value for the thread becomes YES and the HISTORY value becomes NO.

- If `joe` connects from any other host, there is no match. The `INSTRUMENTED` and `HISTORY` values for the thread become `NO`.
- If `sam` connects from any host, the connection matches the third inserted row. The `INSTRUMENTED` value for the thread becomes `NO` and the `HISTORY` value becomes `YES`.
- For any other connection, the row with `HOST` and `USER` set to `'%'` matches. This row now has `ENABLED` and `HISTORY` set to `NO`, so the `INSTRUMENTED` and `HISTORY` values for the thread become `NO`.

Modifications to the `setup_actors` table affect only foreground threads created subsequent to the modification, not existing threads. To affect existing threads, modify the `INSTRUMENTED` and `HISTORY` columns of `threads` table rows.

25.4.7 Pre-Filtering by Consumer

The `setup_consumers` table lists the available consumer types and which are enabled:

```
mysql> SELECT * FROM setup_consumers;
+--------------------------------+---------+
| NAME                           | ENABLED |
+--------------------------------+---------+
| events_stages_current          | NO      |
| events_stages_history          | NO      |
| events_stages_history_long     | NO      |
| events_statements_current      | YES     |
| events_statements_history      | YES     |
| events_statements_history_long | NO      |
| events_transactions_current    | YES     |
| events_transactions_history    | YES     |
| events_transactions_history_long | NO    |
| events_waits_current           | NO      |
| events_waits_history           | NO      |
| events_waits_history_long      | NO      |
| global_instrumentation         | YES     |
| thread_instrumentation         | YES     |
| statements_digest              | YES     |
+--------------------------------+---------+
```

Modify the `setup_consumers` table to affect pre-filtering at the consumer stage and determine the destinations to which events are sent. To enable or disable a consumer, set its `ENABLED` value to `YES` or `NO`.

Modifications to the `setup_consumers` table affect monitoring immediately.

If you disable a consumer, the server does not spend time maintaining destinations for that consumer. For example, if you do not care about historical event information, disable the history consumers:

```
mysql> UPDATE setup_consumers
       SET ENABLED = 'NO' WHERE NAME LIKE '%history%';
```

The consumer settings in the `setup_consumers` table form a hierarchy from higher levels to lower. The following principles apply:

- Destinations associated with a consumer receive no events unless the Performance Schema checks the consumer and the consumer is enabled.
- A consumer is checked only if all consumers it depends on (if any) are enabled.
- If a consumer is not checked, or is checked but is disabled, other consumers that depend on it are not checked.
- Dependent consumers may have their own dependent consumers.
- If an event would not be sent to any destination, the Performance Schema does not produce it.

The following lists describe the available consumer values. For discussion of several representative consumer configurations and their effect on instrumentation, see Section 25.4.8, “Example Consumer Configurations”.

- Global and Thread Consumers
- Wait Event Consumers
- Stage Event Consumers
- Statement Event Consumers
- Transaction Event Consumers
- Statement Digest Consumer

Global and Thread Consumers

- `global_instrumentation` is the highest level consumer. If `global_instrumentation` is `NO`, it disables global instrumentation. All other settings are lower level and are not checked; it does not matter what they are set to. No global or per thread information is maintained and no individual events are collected in the current-events or event-history tables. If `global_instrumentation` is `YES`, the Performance Schema maintains information for global states and also checks the `thread_instrumentation` consumer.
- `thread_instrumentation` is checked only if `global_instrumentation` is `YES`. Otherwise, if `thread_instrumentation` is `NO`, it disables thread-specific instrumentation and all lower-level settings are ignored. No information is maintained per thread and no individual events are collected in the current-events or event-history tables. If `thread_instrumentation` is `YES`, the Performance Schema maintains thread-specific information and also checks `events_xxx_current` consumers.

Wait Event Consumers

These consumers require both `global_instrumentation` and `thread_instrumentation` to be `YES` or they are not checked. If checked, they act as follows:

- `events_waits_current`, if `NO`, disables collection of individual wait events in the `events_waits_current` table. If `YES`, it enables wait event collection and the Performance Schema checks the `events_waits_history` and `events_waits_history_long` consumers.
- `events_waits_history` is not checked if `event_waits_current` is `NO`. Otherwise, an `events_waits_history` value of `NO` or `YES` disables or enables collection of wait events in the `events_waits_history` table.
- `events_waits_history_long` is not checked if `event_waits_current` is `NO`. Otherwise, an `events_waits_history_long` value of `NO` or `YES` disables or enables collection of wait events in the `events_waits_history_long` table.

Stage Event Consumers

These consumers require both `global_instrumentation` and `thread_instrumentation` to be `YES` or they are not checked. If checked, they act as follows:

- `events_stages_current`, if `NO`, disables collection of individual stage events in the `events_stages_current` table. If `YES`, it enables stage event collection and the Performance Schema checks the `events_stages_history` and `events_stages_history_long` consumers.
- `events_stages_history` is not checked if `event_stages_current` is `NO`. Otherwise, an `events_stages_history` value of `NO` or `YES` disables or enables collection of stage events in the `events_stages_history` table.

- `events_stages_history_long` is not checked if `event_stages_current` is `NO`. Otherwise, an `events_stages_history_long` value of `NO` or `YES` disables or enables collection of stage events in the `events_stages_history_long` table.

Statement Event Consumers

These consumers require both `global_instrumentation` and `thread_instrumentation` to be `YES` or they are not checked. If checked, they act as follows:

- `events_statements_current`, if `NO`, disables collection of individual statement events in the `events_statements_current` table. If `YES`, it enables statement event collection and the Performance Schema checks the `events_statements_history` and `events_statements_history_long` consumers.
- `events_statements_history` is not checked if `events_statements_current` is `NO`. Otherwise, an `events_statements_history` value of `NO` or `YES` disables or enables collection of statement events in the `events_statements_history` table.
- `events_statements_history_long` is not checked if `events_statements_current` is `NO`. Otherwise, an `events_statements_history_long` value of `NO` or `YES` disables or enables collection of statement events in the `events_statements_history_long` table.

Transaction Event Consumers

These consumers require both `global_instrumentation` and `thread_instrumentation` to be `YES` or they are not checked. If checked, they act as follows:

- `events_transactions_current`, if `NO`, disables collection of individual transaction events in the `events_transactions_current` table. If `YES`, it enables transaction event collection and the Performance Schema checks the `events_transactions_history` and `events_transactions_history_long` consumers.
- `events_transactions_history` is not checked if `events_transactions_current` is `NO`. Otherwise, an `events_transactions_history` value of `NO` or `YES` disables or enables collection of transaction events in the `events_transactions_history` table.
- `events_transactions_history_long` is not checked if `events_transactions_current` is `NO`. Otherwise, an `events_transactions_history_long` value of `NO` or `YES` disables or enables collection of transaction events in the `events_transactions_history_long` table.

Statement Digest Consumer

The `statements_digest` consumer requires `global_instrumentation` to be `YES` or it is not checked. There is no dependency on the statement event consumers, so you can obtain statistics per digest without having to collect statistics in `events_statements_current`, which is advantageous in terms of overhead. Conversely, you can get detailed statements in `events_statements_current` without digests (the `DIGEST` and `DIGEST_TEXT` columns will be `NULL`).

For more information about statement digesting, see Section 25.9, “Performance Schema Statement Digests”.

25.4.8 Example Consumer Configurations

The consumer settings in the `setup_consumers` table form a hierarchy from higher levels to lower. The following discussion describes how consumers work, showing specific configurations and their effects as consumer settings are enabled progressively from high to low. The consumer values shown are representative. The general principles described here apply to other consumer values that may be available.

The configuration descriptions occur in order of increasing functionality and overhead. If you do not need the information provided by enabling lower-level settings, disable them and the Performance Schema will execute less code on your behalf and you will have less information to sift through.

The `setup_consumers` table contains the following hierarchy of values:

```
global_instrumentation
 thread_instrumentation
   events_waits_current
     events_waits_history
     events_waits_history_long
   events_stages_current
     events_stages_history
     events_stages_history_long
   events_statements_current
     events_statements_history
     events_statements_history_long
   events_transactions_current
     events_transactions_history
     events_transactions_history_long
 statements_digest
```

Note

In the consumer hierarchy, the consumers for waits, stages, statements, and transactions are all at the same level. This differs from the event nesting hierarchy, for which wait events nest within stage events, which nest within statement events, which nest within transaction events.

If a given consumer setting is `NO`, the Performance Schema disables the instrumentation associated with the consumer and ignores all lower-level settings. If a given setting is `YES`, the Performance Schema enables the instrumentation associated with it and checks the settings at the next lowest level. For a description of the rules for each consumer, see Section 25.4.7, “Pre-Filtering by Consumer”.

For example, if `global_instrumentation` is enabled, `thread_instrumentation` is checked. If `thread_instrumentation` is enabled, the `events_xxx_current` consumers are checked. If of these `events_waits_current` is enabled, `events_waits_history` and `events_waits_history_long` are checked.

Each of the following configuration descriptions indicates which setup elements the Performance Schema checks and which output tables it maintains (that is, for which tables it collects information).

- No Instrumentation
- Global Instrumentation Only
- Global and Thread Instrumentation Only
- Global, Thread, and Current-Event Instrumentation
- Global, Thread, Current-Event, and Event-History instrumentation

No Instrumentation

Server configuration state:

```
mysql> SELECT * FROM setup_consumers;
+---------------------------+---------+
| NAME                      | ENABLED |
+---------------------------+---------+
| global_instrumentation    | NO      |
...
+---------------------------+---------+
```

In this configuration, nothing is instrumented.

Setup elements checked:

- Table `setup_consumers`, consumer `global_instrumentation`

Output tables maintained:

- None

Global Instrumentation Only

Server configuration state:

```
mysql> SELECT * FROM setup_consumers;
+---------------------------+---------+
| NAME                      | ENABLED |
+---------------------------+---------+
| global_instrumentation    | YES     |
| thread_instrumentation    | NO      |
...
+---------------------------+---------+
```

In this configuration, instrumentation is maintained only for global states. Per-thread instrumentation is disabled.

Additional setup elements checked, relative to the preceding configuration:

- Table `setup_consumers`, consumer `thread_instrumentation`
- Table `setup_instruments`
- Table `setup_objects`
- Table `setup_timers`

Additional output tables maintained, relative to the preceding configuration:

- `mutex_instances`
- `rwlock_instances`
- `cond_instances`
- `file_instances`
- `users`
- `hosts`
- `accounts`
- `socket_summary_by_event_name`
- `file_summary_by_instance`
- `file_summary_by_event_name`
- `objects_summary_global_by_type`
- `memory_summary_global_by_event_name`
- `table_lock_waits_summary_by_table`
- `table_io_waits_summary_by_index_usage`
- `table_io_waits_summary_by_table`
- `events_waits_summary_by_instance`

- events_waits_summary_global_by_event_name
- events_stages_summary_global_by_event_name
- events_statements_summary_global_by_event_name
- events_transactions_summary_global_by_event_name

Global and Thread Instrumentation Only

Server configuration state:

```
mysql> SELECT * FROM setup_consumers;
+----------------------------------+---------+
| NAME                             | ENABLED |
+----------------------------------+---------+
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| events_waits_current             | NO      |
...
| events_stages_current            | NO      |
...
| events_statements_current        | NO      |
...
| events_transactions_current      | NO      |
...
+----------------------------------+---------+
```

In this configuration, instrumentation is maintained globally and per thread. No individual events are collected in the current-events or event-history tables.

Additional setup elements checked, relative to the preceding configuration:

- Table setup_consumers, consumers events_*xxx*_current, where *xxx* is waits, stages, statements, transactions
- Table setup_actors
- Column threads.instrumented

Additional output tables maintained, relative to the preceding configuration:

- events_*xxx*_summary_by_*yyy*_by_event_name, where *xxx* is waits, stages, statements, transactions; and *yyy* is thread, user, host, account

Global, Thread, and Current-Event Instrumentation

Server configuration state:

```
mysql> SELECT * FROM setup_consumers;
+----------------------------------+---------+
| NAME                             | ENABLED |
+----------------------------------+---------+
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| events_waits_current             | YES     |
| events_waits_history             | NO      |
| events_waits_history_long        | NO      |
| events_stages_current            | YES     |
| events_stages_history            | NO      |
| events_stages_history_long       | NO      |
| events_statements_current        | YES     |
| events_statements_history        | YES     |
| events_statements_history_long   | NO      |
| events_transactions_current      | YES     |
| events_transactions_history      | YES     |
```

```
| events_transactions_history_long | NO      |
...
+----------------------------------+---------+
```

In this configuration, instrumentation is maintained globally and per thread. Individual events are collected in the current-events table, but not in the event-history tables.

Additional setup elements checked, relative to the preceding configuration:

- Consumers `events_xxx_history`, where *xxx* is `waits`, `stages`, `statements`, `transactions`
- Consumers `events_xxx_history_long`, where *xxx* is `waits`, `stages`, `statements`, `transactions`

Additional output tables maintained, relative to the preceding configuration:

- `events_xxx_current`, where *xxx* is `waits`, `stages`, `statements`, `transactions`

Global, Thread, Current-Event, and Event-History instrumentation

The preceding configuration collects no event history because the `events_xxx_history` and `events_xxx_history_long` consumers are disabled. Those consumers can be enabled separately or together to collect event history per thread, globally, or both.

This configuration collects event history per thread, but not globally:

```
mysql> SELECT * FROM setup_consumers;
+----------------------------------+---------+
| NAME                             | ENABLED |
+----------------------------------+---------+
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| events_waits_current             | YES     |
| events_waits_history             | YES     |
| events_waits_history_long        | NO      |
| events_stages_current            | YES     |
| events_stages_history            | YES     |
| events_stages_history_long       | NO      |
| events_statements_current        | YES     |
| events_statements_history        | YES     |
| events_statements_history_long   | NO      |
| events_transactions_current      | YES     |
| events_transactions_history      | YES     |
| events_transactions_history_long | NO      |
...
+----------------------------------+---------+
```

Event-history tables maintained for this configuration:

- `events_xxx_history`, where *xxx* is `waits`, `stages`, `statements`, `transactions`

This configuration collects event history globally, but not per thread:

```
mysql> SELECT * FROM setup_consumers;
+----------------------------------+---------+
| NAME                             | ENABLED |
+----------------------------------+---------+
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| events_waits_current             | YES     |
| events_waits_history             | NO      |
| events_waits_history_long        | YES     |
| events_stages_current            | YES     |
| events_stages_history            | NO      |
| events_stages_history_long       | YES     |
| events_statements_current        | YES     |
| events_statements_history        | YES     |
```

```
| events_statements_history_long   | YES     |
| events_transactions_current      | YES     |
| events_transactions_history      | YES     |
| events_transactions_history_long | YES     |
...
+----------------------------------+---------+
```

Event-history tables maintained for this configuration:

- `events_xxx_history_long`, where *`xxx`* is `waits`, `stages`, `statements`, `transactions`

This configuration collects event history per thread and globally:

```
mysql> SELECT * FROM setup_consumers;
+----------------------------------+---------+
| NAME                             | ENABLED |
+----------------------------------+---------+
| global_instrumentation           | YES     |
| thread_instrumentation           | YES     |
| events_waits_current             | YES     |
| events_waits_history             | YES     |
| events_waits_history_long        | YES     |
| events_stages_current            | YES     |
| events_stages_history            | YES     |
| events_stages_history_long       | YES     |
| events_statements_current        | YES     |
| events_statements_history        | YES     |
| events_statements_history_long   | YES     |
| events_transactions_current      | YES     |
| events_transactions_history      | YES     |
| events_transactions_history_long | YES     |
...
+----------------------------------+---------+
```

Event-history tables maintained for this configuration:

- `events_xxx_history`, where *`xxx`* is `waits`, `stages`, `statements`, `transactions`
- `events_xxx_history_long`, where *`xxx`* is `waits`, `stages`, `statements`, `transactions`

25.4.9 Naming Instruments or Consumers for Filtering Operations

Names given for filtering operations can be as specific or general as required. To indicate a single instrument or consumer, specify its name in full:

```
mysql> UPDATE setup_instruments
       SET ENABLED = 'NO'
       WHERE NAME = 'wait/synch/mutex/myisammrg/MYRG_INFO::mutex';

mysql> UPDATE setup_consumers
       SET ENABLED = 'NO' WHERE NAME = 'events_waits_current';
```

To specify a group of instruments or consumers, use a pattern that matches the group members:

```
mysql> UPDATE setup_instruments
       SET ENABLED = 'NO'
       WHERE NAME LIKE 'wait/synch/mutex/%';

mysql> UPDATE setup_consumers
       SET ENABLED = 'NO' WHERE NAME LIKE '%history%';
```

If you use a pattern, it should be chosen so that it matches all the items of interest and no others. For example, to select all file I/O instruments, it is better to use a pattern that includes the entire instrument name prefix:

```
... WHERE NAME LIKE 'wait/io/file/%';
```

A pattern of `'%/file/%'` will match other instruments that have a component of `'/file/'` anywhere in the name. Even less suitable is the pattern `'%file%'` because it will match instruments with `'file'` anywhere in the name, such as `wait/synch/mutex/sql/LOCK_des_key_file`.

To check which instrument or consumer names a pattern matches, perform a simple test:

```
mysql> SELECT NAME FROM setup_instruments WHERE NAME LIKE 'pattern';

mysql> SELECT NAME FROM setup_consumers WHERE NAME LIKE 'pattern';
```

For information about the types of names that are supported, see Section 25.6, “Performance Schema Instrument Naming Conventions”.

25.4.10 Determining What Is Instrumented

It is always possible to determine what instruments the Performance Schema includes by checking the `setup_instruments` table. For example, to see what file-related events are instrumented for the `InnoDB` storage engine, use this query:

```
mysql> SELECT * FROM setup_instruments WHERE NAME LIKE 'wait/io/file/innodb/%';
+----------------------------------------+---------+-------+
| NAME                                   | ENABLED | TIMED |
+----------------------------------------+---------+-------+
| wait/io/file/innodb/innodb_data_file   | YES     | YES   |
| wait/io/file/innodb/innodb_log_file    | YES     | YES   |
| wait/io/file/innodb/innodb_temp_file   | YES     | YES   |
+----------------------------------------+---------+-------+
```

An exhaustive description of precisely what is instrumented is not given in this documentation, for several reasons:

- What is instrumented is the server code. Changes to this code occur often, which also affects the set of instruments.
- It is not practical to list all the instruments because there are hundreds of them.
- As described earlier, it is possible to find out by querying the `setup_instruments` table. This information is always up to date for your version of MySQL, also includes instrumentation for instrumented plugins you might have installed that are not part of the core server, and can be used by automated tools.

25.5 Performance Schema Queries

Pre-filtering limits which event information is collected and is independent of any particular user. By contrast, post-filtering is performed by individual users through the use of queries with appropriate `WHERE` clauses that restrict what event information to select from the events available after pre-filtering has been applied.

In Section 25.4.3, “Event Pre-Filtering”, an example showed how to pre-filter for file instruments. If the event tables contain both file and nonfile information, post-filtering is another way to see information only for file events. Add a `WHERE` clause to queries to restrict event selection appropriately:

```
mysql> SELECT THREAD_ID, NUMBER_OF_BYTES
       FROM events_waits_history
       WHERE EVENT_NAME LIKE 'wait/io/file/%'
       AND NUMBER_OF_BYTES IS NOT NULL;
+-----------+-----------------+
| THREAD_ID | NUMBER_OF_BYTES |
+-----------+-----------------+
|        11 |              66 |
```

```
|        11 |              47 |
|        11 |             139 |
|         5 |              24 |
|         5 |             834 |
+-----------+-----------------+
```

Most Performance Schema tables have indexes, which gives the optimizer access to execution plans other than full table scans. These indexes also improve performance for related objects, such as `sys` schema views that use those tables. For more information, see Section 8.2.4, “Optimizing Performance Schema Queries”.

25.6 Performance Schema Instrument Naming Conventions

An instrument name consists of a sequence of components separated by `'/'` characters. Example names:

```
wait/io/file/myisam/log
wait/io/file/mysys/charset
wait/lock/table/sql/handler
wait/synch/cond/mysys/COND_alarm
wait/synch/cond/sql/BINLOG::update_cond
wait/synch/mutex/mysys/BITMAP_mutex
wait/synch/mutex/sql/LOCK_delete
wait/synch/rwlock/sql/Query_cache_query::lock
stage/sql/closing tables
stage/sql/Sorting result
statement/com/Execute
statement/com/Query
statement/sql/create_table
statement/sql/lock_tables
errors
```

The instrument name space has a tree-like structure. The components of an instrument name from left to right provide a progression from more general to more specific. The number of components a name has depends on the type of instrument.

The interpretation of a given component in a name depends on the components to the left of it. For example, `myisam` appears in both of the following names, but `myisam` in the first name is related to file I/O, whereas in the second it is related to a synchronization instrument:

```
wait/io/file/myisam/log
wait/synch/cond/myisam/MI_SORT_INFO::cond
```

Instrument names consist of a prefix with a structure defined by the Performance Schema implementation and a suffix defined by the developer implementing the instrument code. The top-level component of an instrument prefix indicates the type of instrument. This component also determines which event timer in the `setup_timers` table applies to the instrument. For the prefix part of instrument names, the top level indicates the type of instrument.

The suffix part of instrument names comes from the code for the instruments themselves. Suffixes may include levels such as these:

- A name for the major component (a server module such as `myisam`, `innodb`, `mysys`, or `sql`) or a plugin name.
- The name of a variable in the code, in the form *`XXX`* (a global variable) or *`CCC`*`::`*`MMM`* (a member *`MMM`* in class *`CCC`*). Examples: `COND_thread_cache`, `THR_LOCK_myisam`, `BINLOG::LOCK_index`.
- Top-Level Instrument Components
- Idle Instrument Components
- Error Instrument Components

- Memory Instrument Components
- Stage Instrument Components
- Statement Instrument Components
- Wait Instrument Components

Top-Level Instrument Components

- `idle`: An instrumented idle event. This instrument has no further components.
- `error`: An instrumented error event. This instrument has no further components.
- `memory`: An instrumented memory event.
- `stage`: An instrumented stage event.
- `statement`: An instrumented statement event.
- `transaction`: An instrumented transaction event. This instrument has no further components.
- `wait`: An instrumented wait event.

Idle Instrument Components

The `idle` instrument is used for idle events, which The Performance Schema generates as discussed in the description of the `socket_instances.STATE` column in Section 25.11.3.5, “The socket_instances Table”.

Error Instrument Components

The `error` instrument indicates whether to collect information for server errors and warnings. This instrument is enabled by default. The `TIMED` column for the `error` row in the `setup_instruments` table is inapplicable because timing information is not collected.

Memory Instrument Components

Memory instrumentation is enabled by default. (Prior to 8.0.2, most memory instrumentation is disabled by default.) Memory instrumentation can be enabled or disabled at startup, or dynamically at runtime by updating the `ENABLED` column of the relevant instruments in the `setup_instruments` table. Memory instruments have names of the form `memory/`*`code_area`*`/`*`instrument_name`* where *`code_area`* is a value such as `sql` or `myisam`, and *`instrument_name`* is the instrument detail.

Instruments named with the prefix `memory/performance_schema/` expose how much memory is allocated for internal buffers in the Performance Schema. The `memory/performance_schema/` instruments are built in, always enabled, and cannot be disabled at startup or runtime. Built-in memory instruments are displayed only in the `memory_summary_global_by_event_name` table. For more information, see Section 25.16, “The Performance Schema Memory-Allocation Model”.

Stage Instrument Components

Stage instruments have names of the form `stage/`*`code_area`*`/`*`stage_name`*, where *`code_area`* is a value such as `sql` or `myisam`, and *`stage_name`* indicates the stage of statement processing, such as `Sorting result` or `Sending data`. Stages correspond to the thread states displayed by `SHOW PROCESSLIST` or that are visible in the `INFORMATION_SCHEMA.PROCESSLIST` table.

Statement Instrument Components

- `statement/abstract/*`: An abstract instrument for statement operations. Abstract instruments are used during the early stages of statement classification before the exact statement type is

known, then changed to a more specific statement instrument when the type is known. For a description of this process, see Section 25.11.6, “Performance Schema Statement Event Tables”.

- `statement/com`: An instrumented command operation. These have names corresponding to `COM_xxx` operations (see the `mysql_com.h` header file and `sql/sql_parse.cc`. For example, the `statement/com/Connect` and `statement/com/Init DB` instruments correspond to the `COM_CONNECT` and `COM_INIT_DB` commands.
- `statement/scheduler/event`: A single instrument to track all events executed by the Event Scheduler. This instrument comes into play when a scheduled event begins executing.
- `statement/sp`: An instrumented internal instruction executed by a stored program. For example, the `statement/sp/cfetch` and `statement/sp/freturn` instruments are used cursor fetch and function return instructions.
- `statement/sql`: An instrumented SQL statement operation. For example, the `statement/sql/create_db` and `statement/sql/select` instruments are used for `CREATE DATABASE` and `SELECT` statements.

Wait Instrument Components

- `wait/io`

  An instrumented I/O operation.

  - `wait/io/file`

    An instrumented file I/O operation. For files, the wait is the time waiting for the file operation to complete (for example, a call to `fwrite()`). Due to caching, the physical file I/O on the disk might not happen within this call.

  - `wait/io/socket`

    An instrumented socket operation. Socket instruments have names of the form `wait/io/socket/sql/socket_type`. The server has a listening socket for each network protocol that it supports. The instruments associated with listening sockets for TCP/IP or Unix socket file connections have a *`socket_type`* value of `server_tcpip_socket` or `server_unix_socket`, respectively. When a listening socket detects a connection, the server transfers the connection to a new socket managed by a separate thread. The instrument for the new connection thread has a *`socket_type`* value of `client_connection`.

  - `wait/io/table`

    An instrumented table I/O operation. These include row-level accesses to persistent base tables or temporary tables. Operations that affect rows are fetch, insert, update, and delete. For a view, waits are associated with base tables referenced by the view.

    Unlike most waits, a table I/O wait can include other waits. For example, table I/O might include file I/O or memory operations. Thus, `events_waits_current` for a table I/O wait usually has two rows. For more information, see Section 25.8, “Performance Schema Atom and Molecule Events”.

    Some row operations might cause multiple table I/O waits. For example, an insert might activate a trigger that causes an update.

- `wait/lock`

  An instrumented lock operation.

  - `wait/lock/table`

    An instrumented table lock operation.

  - `wait/lock/metadata/sql/mdl`

    An instrumented metadata lock operation.

- `wait/synch`

  An instrumented synchronization object. For synchronization objects, the `TIMER_WAIT` time includes the amount of time blocked while attempting to acquire a lock on the object, if any.

  - `wait/synch/cond`

    A condition is used by one thread to signal to other threads that something they were waiting for has happened. If a single thread was waiting for a condition, it can wake up and proceed with its execution. If several threads were waiting, they can all wake up and compete for the resource for which they were waiting.

  - `wait/synch/mutex`

    A mutual exclusion object used to permit access to a resource (such as a section of executable code) while preventing other threads from accessing the resource.

  - `wait/synch/rwlock`

    A read/write lock object used to lock a specific variable for access while preventing its use by other threads. A shared read lock can be acquired simultaneously by multiple threads. An exclusive write lock can be acquired by only one thread at a time.

  - `wait/synch/sxlock`

    A shared-exclusive (SX) lock is a type of rwlock lock object that provides write access to a common resource while permitting inconsistent reads by other threads. `sxlocks` optimize concurrency and improve scalability for read-write workloads.

25.7 Performance Schema Status Monitoring

There are several status variables associated with the Performance Schema:

```
mysql> SHOW STATUS LIKE 'perf%';
+-----------------------------------------------+-------+
| Variable_name                                 | Value |
+-----------------------------------------------+-------+
| Performance_schema_accounts_lost              | 0     |
| Performance_schema_cond_classes_lost          | 0     |
| Performance_schema_cond_instances_lost        | 0     |
| Performance_schema_digest_lost                | 0     |
| Performance_schema_file_classes_lost          | 0     |
| Performance_schema_file_handles_lost          | 0     |
| Performance_schema_file_instances_lost        | 0     |
| Performance_schema_hosts_lost                 | 0     |
| Performance_schema_locker_lost                | 0     |
| Performance_schema_memory_classes_lost        | 0     |
| Performance_schema_metadata_lock_lost         | 0     |
| Performance_schema_mutex_classes_lost         | 0     |
| Performance_schema_mutex_instances_lost       | 0     |
| Performance_schema_nested_statement_lost      | 0     |
| Performance_schema_program_lost               | 0     |
| Performance_schema_rwlock_classes_lost        | 0     |
| Performance_schema_rwlock_instances_lost      | 0     |
| Performance_schema_session_connect_attrs_lost | 0     |
| Performance_schema_socket_classes_lost        | 0     |
| Performance_schema_socket_instances_lost      | 0     |
| Performance_schema_stage_classes_lost         | 0     |
| Performance_schema_statement_classes_lost     | 0     |
| Performance_schema_table_handles_lost         | 0     |
```

```
| Performance_schema_table_instances_lost       | 0     |
| Performance_schema_thread_classes_lost        | 0     |
| Performance_schema_thread_instances_lost      | 0     |
| Performance_schema_users_lost                 | 0     |
+-----------------------------------------------+-------+
```

The Performance Schema status variables provide information about instrumentation that could not be loaded or created due to memory constraints. Names for these variables have several forms:

- `Performance_schema_xxx_classes_lost` indicates how many instruments of type *`xxx`* could not be loaded.
- `Performance_schema_xxx_instances_lost` indicates how many instances of object type *`xxx`* could not be created.
- `Performance_schema_xxx_handles_lost` indicates how many instances of object type *`xxx`* could not be opened.
- `Performance_schema_locker_lost` indicates how many events are “lost” or not recorded.

For example, if a mutex is instrumented in the server source but the server cannot allocate memory for the instrumentation at runtime, it increments `Performance_schema_mutex_classes_lost`. The mutex still functions as a synchronization object (that is, the server continues to function normally), but performance data for it will not be collected. If the instrument can be allocated, it can be used for initializing instrumented mutex instances. For a singleton mutex such as a global mutex, there will be only one instance. Other mutexes have an instance per connection, or per page in various caches and data buffers, so the number of instances varies over time. Increasing the maximum number of connections or the maximum size of some buffers will increase the maximum number of instances that might be allocated at once. If the server cannot create a given instrumented mutex instance, it increments `Performance_schema_mutex_instances_lost`.

Suppose that the following conditions hold:

- The server was started with the `--performance_schema_max_mutex_classes=200` option and thus has room for 200 mutex instruments.
- 150 mutex instruments have been loaded already.
- The plugin named `plugin_a` contains 40 mutex instruments.
- The plugin named `plugin_b` contains 20 mutex instruments.

The server allocates mutex instruments for the plugins depending on how many they need and how many are available, as illustrated by the following sequence of statements:

```
INSTALL PLUGIN plugin_a
```

The server now has 150+40 = 190 mutex instruments.

```
UNINSTALL PLUGIN plugin_a;
```

The server still has 190 instruments. All the historical data generated by the plugin code is still available, but new events for the instruments are not collected.

```
INSTALL PLUGIN plugin_a;
```

The server detects that the 40 instruments are already defined, so no new instruments are created, and previously assigned internal memory buffers are reused. The server still has 190 instruments.

```
INSTALL PLUGIN plugin_b;
```

The server has room for 200-190 = 10 instruments (in this case, mutex classes), and sees that the plugin contains 20 new instruments. 10 instruments are loaded, and 10 are discarded or “lost.” The `Performance_schema_mutex_classes_lost` indicates the number of instruments (mutex classes) lost:

```
mysql> SHOW STATUS LIKE "perf%mutex_classes_lost";
+---------------------------------------+-------+
| Variable_name                         | Value |
+---------------------------------------+-------+
| Performance_schema_mutex_classes_lost | 10    |
+---------------------------------------+-------+
1 row in set (0.10 sec)
```

The instrumentation still works and collects (partial) data for `plugin_b`.

When the server cannot create a mutex instrument, these results occur:

- No row for the instrument is inserted into the `setup_instruments` table.
- `Performance_schema_mutex_classes_lost` increases by 1.
- `Performance_schema_mutex_instances_lost` does not change. (When the mutex instrument is not created, it cannot be used to create instrumented mutex instances later.)

The pattern just described applies to all types of instruments, not just mutexes.

A value of `Performance_schema_mutex_classes_lost` greater than 0 can happen in two cases:

- To save a few bytes of memory, you start the server with `--performance_schema_max_mutex_classes=N`, where *`N`* is less than the default value. The default value is chosen to be sufficient to load all the plugins provided in the MySQL distribution, but this can be reduced if some plugins are never loaded. For example, you might choose not to load some of the storage engines in the distribution.
- You load a third-party plugin that is instrumented for the Performance Schema but do not allow for the plugin's instrumentation memory requirements when you start the server. Because it comes from a third party, the instrument memory consumption of this engine is not accounted for in the default value chosen for `performance_schema_max_mutex_classes`.

  If the server has insufficient resources for the plugin's instruments and you do not explicitly allocate more using `--performance_schema_max_mutex_classes=N`, loading the plugin leads to starvation of instruments.

If the value chosen for `performance_schema_max_mutex_classes` is too small, no error is reported in the error log and there is no failure at runtime. However, the content of the tables in the `performance_schema` database will miss events. The `Performance_schema_mutex_classes_lost` status variable is the only visible sign to indicate that some events were dropped internally due to failure to create instruments.

If an instrument is not lost, it is known to the Performance Schema, and is used when instrumenting instances. For example, `wait/synch/mutex/sql/LOCK_delete` is the name of a mutex instrument in the `setup_instruments` table. This single instrument is used when creating a mutex in the code (in `THD::LOCK_delete`) however many instances of the mutex are needed as the server runs. In this case, `LOCK_delete` is a mutex that is per connection (`THD`), so if a server has 1000 connections, there are 1000 threads, and 1000 instrumented `LOCK_delete` mutex instances (`THD::LOCK_delete`).

If the server does not have room for all these 1000 instrumented mutexes (instances), some mutexes are created with instrumentation, and some are created without instrumentation. If the server can create only 800 instances, 200 instances are lost. The server continues to run, but increments `Performance_schema_mutex_instances_lost` by 200 to indicate that instances could not be created.

A value of `Performance_schema_mutex_instances_lost` greater than 0 can happen when the code initializes more mutexes at runtime than were allocated for `--performance_schema_max_mutex_instances=N`.

The bottom line is that if `SHOW STATUS LIKE 'perf%'` says that nothing was lost (all values are zero), the Performance Schema data is accurate and can be relied upon. If something was lost, the data is incomplete, and the Performance Schema could not record everything given the insufficient amount of memory it was given to use. In this case, the specific `Performance_schema_xxx_lost` variable indicates the problem area.

It might be appropriate in some cases to cause deliberate instrument starvation. For example, if you do not care about performance data for file I/O, you can start the server with all Performance Schema parameters related to file I/O set to 0. No memory will be allocated for file-related classes, instances, or handles, and all file events will be lost.

Use `SHOW ENGINE PERFORMANCE_SCHEMA STATUS` to inspect the internal operation of the Performance Schema code:

```
mysql> SHOW ENGINE PERFORMANCE_SCHEMA STATUS\G
...
*************************** 3. row ***************************
  Type: performance_schema
  Name: events_waits_history.size
Status: 76
*************************** 4. row ***************************
  Type: performance_schema
  Name: events_waits_history.count
Status: 10000
*************************** 5. row ***************************
  Type: performance_schema
  Name: events_waits_history.memory
Status: 760000
...
*************************** 57. row ***************************
  Type: performance_schema
  Name: performance_schema.memory
Status: 26459600
...
```

This statement is intended to help the DBA understand the effects that different Performance Schema options have on memory requirements. For a description of the field meanings, see Section 13.7.5.15, “SHOW ENGINE Syntax”.

25.8 Performance Schema Atom and Molecule Events

For a table I/O event, there are usually two rows in `events_waits_current`, not one. For example, a row fetch might result in rows like this:

```
Row# EVENT_NAME                TIMER_START TIMER_END
---- ----------                ----------- ---------
   1 wait/io/file/myisam/dfile       10001 10002
   2 wait/io/table/sql/handler       10000 NULL
```

The row fetch causes a file read. In the example, the table I/O fetch event started before the file I/O event but has not finished (its `TIMER_END` value is `NULL`). The file I/O event is “nested” within the table I/O event.

This occurs because, unlike other “atomic” wait events such as for mutexes or file I/O, table I/O events are “molecular” and include (overlap with) other events. In `events_waits_current`, the table I/O event usually has two rows:

- One row for the most recent table I/O wait event

- One row for the most recent wait event of any kind

Usually, but not always, the “of any kind” wait event differs from the table I/O event. As each subsidiary event completes, it disappears from `events_waits_current`. At this point, and until the next subsidiary event begins, the table I/O wait is also the most recent wait of any kind.

25.9 Performance Schema Statement Digests

The MySQL server is capable of maintaining statement digest information. The digesting process converts each SQL statement to normalized form and computes a MD5 hash value from the normalized result. Normalization permits statements that are similar to be grouped and summarized to expose information about the types of statements the server is executing and how often they occur. This section describes how statement digesting occurs and how it can be useful.

Digesting occurs in the parser regardless of whether the Performance Schema is available, so that other server components such as query rewrite plugins have access to statement digests.

- Statement Digest General Concepts
- Statement Digests in the Performance Schema
- Statement Digest Memory Use

Statement Digest General Concepts

When the parser receives an SQL statement, it computes a statement digest if that digest is needed, which is true if any of the following conditions are true:

- Performance Schema digest instrumentation is enabled
- A Query Rewrite Plugin is enabled

The `max_digest_length` system variable value determines the maximum number of bytes available for computing normalized statement digests. Once that amount of space is used during digest computation, truncation occurs: no further tokens from a parsed statement are collected or figure into its digest value. Statements that differ only after that many bytes of parsed tokens produce the same normalized statement digest and are considered identical if compared or if aggregated for digest statistics.

After the normalized statement has been computed, an MD5 hash value is computed from it. In addition:

- If any Query Rewrite Plugin is enabled, it is called and the digest as computed is available to it.
- If the Performance Schema has digest instrumentation enabled, it makes a copy of the normalized statement digest, allocating a maximum of `performance_schema_max_digest_length` bytes for it. Consequently, if `performance_schema_max_digest_length` is less than `max_digest_length`, the copy is truncated relative to the original. The copy of the normalized statement is stored in the appropriate Performance Schema tables, along with the MD5 hash value computed from the original normalized statement. (If the Performance Schema truncates its copy of the normalized statement relative to the original, it does not recompute the MD5 hash value.)

Statement normalization transforms the statement text to a more standardized digest string representation that preserves the general statement structure while removing information not essential to the structure:

- Object identifiers such as database and table names are preserved.
- Literal values are converted to parameter markers. A normalized statement does not retain information such as names, passwords, dates, and so forth.
- Comments are removed and whitespace is adjusted.

Consider these statements:

```
SELECT * FROM orders WHERE customer_id=10 AND quantity>20
SELECT * FROM orders WHERE customer_id = 20 AND quantity > 100
```

To normalize these statements, the parser replaces data values by ? and adjusts whitespace. Both statements yield the same normalized form and thus are considered “the same”:

```
SELECT * FROM orders WHERE customer_id = ? AND quantity > ?
```

The normalized statement contains less information but is still representative of the original statement. Other similar statements that have different data values have the same normalized form.

Now consider these statements:

```
SELECT * FROM customers WHERE customer_id = 1000
SELECT * FROM orders WHERE customer_id = 1000
```

In this case, the normalized statements differ because the object identifiers differ:

```
SELECT * FROM customers WHERE customer_id = ?
SELECT * FROM orders WHERE customer_id = ?
```

If normalization produces a statement that exceeds the space available in the digest buffer (as determined by `max_digest_length`), truncation occurs and the text ends with “...”. Long normalized statements that differ only in the part that occurs following the “...” are considered the same. Consider these statements:

```
SELECT * FROM mytable WHERE cola = 10 AND colb = 20
SELECT * FROM mytable WHERE cola = 10 AND colc = 20
```

If the cutoff happens to be right after the `AND`, both statements have this normalized form:

```
SELECT * FROM mytable WHERE cola = ? AND ...
```

In this case, the difference in the second column name is lost and both statements are considered the same.

Statement Digests in the Performance Schema

In the Performance Schema, statement digesting involves these components:

- A `statements_digest` consumer in the `setup_consumers` table controls whether the Performance Schema maintains digest information. See Statement Digest Consumer.
- The statement event tables (`events_statements_current`, `events_statements_history`, and `events_statements_history_long`) have columns for storing normalized statement digests and the corresponding digest MD5 hash values:
  - `DIGEST_TEXT` is the text of the normalized statement digest. This is a copy of the original normalized statement that was computed to a maximum of `max_digest_length` bytes, further truncated as necessary to `performance_schema_max_digest_length` bytes.
  - `DIGEST` is the digest MD5 hash value computed from the original normalized statement.

  The statement event tables also have an `SQL_TEXT` column that contains the original SQL statement. The maximum space available for statement display is 1024 bytes by default. To change this value, set the `performance_schema_max_sql_text_length` system variable at server startup.

See Section 25.11.6, “Performance Schema Statement Event Tables”.

- The `events_statements_summary_by_digest` summary table provides aggregated statement digest information. This table aggregates information for statements that have the same `SCHEMA_NAME` and `DIGEST` values. The Performance Schema uses MD5 hash values for aggregation because they are fast to compute and have a favorable statistical distribution that minimizes collisions. See Section 25.11.15.3, “Statement Summary Tables”.

The `performance_schema_max_digest_length` system variable determines the maximum number of bytes available for digest value storage in the Performance Schema. However, the display length of statement digests may be longer than the available buffer size due to internal encoding of statement components such as keywords and literal values. Consequently, values selected from the `DIGEST_TEXT` column of statement event tables may appear to exceed the `performance_schema_max_digest_length` value.

The `events_statements_summary_by_digest` summary table provides a profile of the statements executed by the server. It shows what kinds of statements an application is executing and how often. An application developer can use this information together with other information in the table to assess the application's performance characteristics. For example, table columns that show wait times, lock times, or index use may highlight types of queries that are inefficient. This gives the developer insight into which parts of the application need attention.

The `events_statements_summary_by_digest` summary table has a fixed size. By default the Performance Schema estimates the size to use at startup. To specify the table size explicitly, set the `performance_schema_digests_size` system variable at server startup. If the table becomes full, the Performance Schema groups statements that have `SCHEMA_NAME` and `DIGEST` values not matching existing values in the table in a special row with `SCHEMA_NAME` and `DIGEST` set to `NULL`. This permits all statements to be counted. However, if the special row accounts for a significant percentage of the statements executed, it might be desirable to increase the summary table size by increasing `performance_schema_digests_size`.

Statement Digest Memory Use

For applications that generate very long statements that differ only at the end, increasing `max_digest_length` enables computation of digests that distinguish statements that would otherwise aggregate to the same digest. Conversely, decreasing `max_digest_length` causes the server to devote less memory to digest storage but increases the likelihood of longer statements aggregating to the same digest. Administrators should keep in mind that larger values result in correspondingly increased memory requirements, particularly for workloads that involve large numbers of simultaneous sessions (the server allocates `max_digest_length` bytes per session).

As described previously, normalized statement digests as computed by the parser are constrained to a maximum of `max_digest_length` bytes, whereas normalized statement digests stored in the Performance Schema use `performance_schema_max_digest_length` bytes. The following memory-use considerations apply regarding the relative values of `max_digest_length` and `performance_schema_max_digest_length`:

- If `max_digest_length` is less than `performance_schema_max_digest_length`:
  - Server components other than the Performance Schema use normalized statement digests that take up to `max_digest_length` bytes.
  - The Performance Schema does not further truncate normalized statement digests that it stores, but allocates more memory than `max_digest_length` bytes per digest, which is unnecessary.
- If `max_digest_length` equals `performance_schema_max_digest_length`:
  - Server components other than the Performance Schema use normalized statement digests that take up to `max_digest_length` bytes.

  - The Performance Schema does not further truncate normalized statement digests that it stores, and allocates the same amount of memory as `max_digest_length` bytes per digest.

- If `max_digest_length` is greater than `performance_schema_max_digest_length`:
  - Server components other than the Performance Schema use normalized statement digests that take up to `max_digest_length` bytes.
  - The Performance Schema further truncates normalized statement digests that it stores, and allocates less memory than `max_digest_length` bytes per digest.

Because the Performance Schema statement event tables might store many digests, setting `performance_schema_max_digest_length` smaller than `max_digest_length` enables administrators to balance these factors:

- The need to have long normalized statement digests available for server components outside the Performance Schema
- Many concurrent sessions, each of which allocates digest-computation memory
- The need to limit memory consumption by the Performance Schema statement event tables when storing many statement digests

To assess the amount of memory used for SQL statement storage and digest computation, use the `SHOW ENGINE PERFORMANCE_SCHEMA STATUS` statement, or monitor these instruments:

```
mysql> SELECT NAME FROM setup_instruments
       WHERE NAME LIKE '%.sqltext';
+-------------------------------------------------------------------+
| NAME                                                              |
+-------------------------------------------------------------------+
| memory/performance_schema/events_statements_history.sqltext       |
| memory/performance_schema/events_statements_current.sqltext       |
| memory/performance_schema/events_statements_history_long.sqltext |
+-------------------------------------------------------------------+

mysql> SELECT NAME FROM setup_instruments
       WHERE NAME LIKE 'memory/performance_schema/%.tokens';
+-----------------------------------------------------------------------+
| NAME                                                                  |
+-----------------------------------------------------------------------+
| memory/performance_schema/events_statements_history.tokens            |
| memory/performance_schema/events_statements_current.tokens            |
| memory/performance_schema/events_statements_summary_by_digest.tokens |
| memory/performance_schema/events_statements_history_long.tokens       |
+-----------------------------------------------------------------------+
```

25.10 Performance Schema General Table Characteristics

The name of the `performance_schema` database is lowercase, as are the names of tables within it. Queries should specify the names in lowercase.

Many tables in the `performance_schema` database are read only and cannot be modified:

```
mysql> TRUNCATE TABLE setup_instruments;
ERROR 1683 (HY000): Invalid performance_schema usage.
```

Some of the setup tables have columns that can be modified to affect Performance Schema operation; some also permit rows to be inserted or deleted. Truncation is permitted to clear collected events, so `TRUNCATE TABLE` can be used on tables containing those kinds of information, such as tables named with a prefix of `events_waits_`.

Summary tables can be truncated with `TRUNCATE TABLE`. Generally, the effect is to reset the summary columns to 0 or `NULL`, not to remove rows. This enables you to clear collected values and restart aggregation. That might be useful, for example, after you have made a runtime configuration change. Exceptions to this truncation behavior are noted in individual summary table sections.

Privileges are as for other databases and tables:

- To retrieve from `performance_schema` tables, you must have the `SELECT` privilege.
- To change those columns that can be modified, you must have the `UPDATE` privilege.
- To truncate tables that can be truncated, you must have the `DROP` privilege.

Because only a limited set of privileges apply to Performance Schema tables, attempts to use `GRANT ALL` as shorthand for granting privileges at the database or table leval fail with an error:

```
mysql> GRANT ALL ON performance_schema.*
       TO 'u1'@'localhost';
ERROR 1044 (42000): Access denied for user 'root'@'localhost'
to database 'performance_schema'
mysql> GRANT ALL ON performance_schema.setup_instruments
       TO 'u2'@'localhost';
ERROR 1044 (42000): Access denied for user 'root'@'localhost'
to database 'performance_schema'
```

Instead, grant exactly the desired privileges:

```
mysql> GRANT SELECT ON performance_schema.*
       TO 'u1'@'localhost';
Query OK, 0 rows affected (0.03 sec)

mysql> GRANT SELECT, UPDATE ON performance_schema.setup_instruments
       TO 'u2'@'localhost';
Query OK, 0 rows affected (0.02 sec)
```

25.11 Performance Schema Table Descriptions

Tables in the `performance_schema` database can be grouped as follows:

- Setup tables. These tables are used to configure and display monitoring characteristics.
- Current events tables. The `events_waits_current` table contains the most recent event for each thread. Other similar tables contain current events at different levels of the event hierarchy: `events_stages_current` for stage events, `events_statements_current` for statement events, and `events_transactions_current` for transaction events.
- History tables. These tables have the same structure as the current events tables, but contain more rows. For example, for wait events, `events_waits_history` table contains the most recent 10 events per thread. `events_waits_history_long` contains the most recent 10,000 events. Other similar tables exist for stage, statement, and transaction histories.

  To change the sizes of the history tables, set the appropriate system variables at server startup. For example, to set the sizes of the wait event history tables, set `performance_schema_events_waits_history_size` and `performance_schema_events_waits_history_long_size`.
- Summary tables. These tables contain information aggregated over groups of events, including those that have been discarded from the history tables.
- Instance tables. These tables document what types of objects are instrumented. An instrumented object, when used by the server, produces an event. These tables provide event names and explanatory notes or status information.
- Miscellaneous tables. These do not fall into any of the other table groups.

25.11.1 Performance Schema Table Index

The following table lists each Performance Schema table and provides a short description of each one.

Table 25.1 Performance Schema Tables

Table Name	Description
accounts	Connection statistics per client account
cond_instances	synchronization object instances
data_lock_waits	Data lock wait relationships
data_locks	Data locks held and requested
events_errors_summary_by_account_by_error	Errors per error code and account
events_errors_summary_by_host_by_error	Errors per error code and host
events_errors_summary_by_thread_by_error	Errors per error code and host
events_errors_summary_by_user_by_error	Errors per error code and host
events_errors_summary_global_by_error	Errors per error code
events_stages_current	Current stage events
events_stages_history	Most recent stage events for each thread
events_stages_history_long	Most recent stage events overall
events_stages_summary_by_account_by_event_name	Stage events per account and event name
events_stages_summary_by_host_by_event_name	Stage events per host name and event name
events_stages_summary_by_thread_by_event_name	Stage waits per thread and event name
events_stages_summary_by_user_by_event_name	Stage events per user name and event name
events_stages_summary_global_by_event_name	Stage waits per event name
events_statements_current	Current statement events
events_statements_histogram_by_digest	Statement histograms per schema and digest value
events_statements_histogram_global	Statement histogram summarized globally
events_statements_history	Most recent statement events for each thread
events_statements_history_long	Most recent statement events overall
events_statements_summary_by_account_by_event_name	Statement events per account and event name
events_statements_summary_by_digest	Statement events per schema and digest value
events_statements_summary_by_host_by_event_name	Statement events per host name and event name
events_statements_summary_by_program	Statement events per stored program
events_statements_summary_by_thread_by_event_name	Statement events per thread and event name
events_statements_summary_by_user_by_event_name	Statement events per user name and event name
events_statements_summary_global_by_event_name	Statement events per event name
events_transactions_current	Current transaction events
events_transactions_history	Most recent transaction events for each thread
events_transactions_history_long	Most recent transaction events overall

Table Name	Description
events_transactions_summary_by_account_by_event_name	Transaction events per account and event name
events_transactions_summary_by_host_by_event_name	Transaction events per host name and event name
events_transactions_summary_by_thread_by_event_name	Transaction events per thread and event name
events_transactions_summary_by_user_by_event_name	Transaction events per user name and event name
events_transactions_summary_global_by_event_name	Transaction events per event name
events_waits_current	Current wait events
events_waits_history	Most recent wait events for each thread
events_waits_history_long	Most recent wait events overall
events_waits_summary_by_account_by_event_name	Wait events per account and event name
events_waits_summary_by_host_by_event_name	Wait events per host name and event name
events_waits_summary_by_instance	Wait events per instance
events_waits_summary_by_thread_by_event_name	Wait events per thread and event name
events_waits_summary_by_user_by_event_name	Wait events per user name and event name
events_waits_summary_global_by_event_name	Wait events per event name
file_instances	File instances
file_summary_by_event_name	File events per event name
file_summary_by_instance	File events per file instance
global_status	Global status variables
global_variables	Global system variables
host_cache	Information from the internal host cache
hosts	Connection statistics per client host name
memory_summary_by_account_by_event_name	Memory operations per account and event name
memory_summary_by_host_by_event_name	Memory operations per host and event name
memory_summary_by_thread_by_event_name	Memory operations per thread and event name
memory_summary_by_user_by_event_name	Memory operations per user and event name
memory_summary_global_by_event_name	Memory operations globally per event name
metadata_locks	Metadata locks and lock requests
mutex_instances	Mutex synchronization object instances
objects_summary_global_by_type	Object summaries
performance_timers	Which event timers are available
persisted_variables	Contents of mysqld-auto.cnf file
prepared_statements_instances	Prepared statement instances and statistics
replication_applier_configuration	Configuration parameters for the transaction applier on the slave
replication_applier_status	Current status of the transaction applier on the slave
replication_applier_status_by_coordinator	SQL or coordinator thread applier status

Table Name	Description
replication_applier_status_by_worker	Worker thread applier status (empty unless slave is multi-threaded)
replication_connection_configuration	Configuration parameters for connecting to the master
replication_connection_status	Current status of the connection to the master
rwlock_instances	Lock synchronization object instances
session_account_connect_attrs	Connection attributes per for the current session
session_connect_attrs	Connection attributes for all sessions
session_status	Status variables for current session
session_variables	System variables for current session
setup_actors	How to initialize monitoring for new foreground threads
setup_consumers	Consumers for which event information can be stored
setup_instruments	Classes of instrumented objects for which events can be collected
setup_objects	Which objects should be monitored
setup_timers	Current event timer
socket_instances	Active connection instances
socket_summary_by_event_name	Socket waits and I/O per event name
socket_summary_by_instance	Socket waits and I/O per instance
status_by_account	Session status variables per account
status_by_host	Session status variables per host name
status_by_thread	Session status variables per session
status_by_user	Session status variables per user name
table_handles	Table locks and lock requests
table_io_waits_summary_by_index_usage	Table I/O waits per index
table_io_waits_summary_by_table	Table I/O waits per table
table_lock_waits_summary_by_table	Table lock waits per table
threads	Information about server threads
user_defined_functions	Registered user-defined functions
user_variables_by_thread	User-defined variables per thread
users	Connection statistics per client user name
variables_by_thread	Session system variables per session
variables_info	How system variables were most recently set

25.11.2 Performance Schema Setup Tables

The setup tables provide information about the current instrumentation and enable the monitoring configuration to be changed. For this reason, some columns in these tables can be changed if you have the UPDATE privilege.

The use of tables rather than individual variables for setup information provides a high degree of flexibility in modifying Performance Schema configuration. For example, you can use a single statement with standard SQL syntax to make multiple simultaneous configuration changes.

These setup tables are available:

- `setup_actors`: How to initialize monitoring for new foreground threads
- `setup_consumers`: The destinations to which event information can be sent and stored
- `setup_instruments`: The classes of instrumented objects for which events can be collected
- `setup_objects`: Which objects should be monitored
- `setup_timers`: The current event timer

25.11.2.1 The setup_actors Table

The `setup_actors` table contains information that determines whether to enable monitoring and historical event logging for new foreground server threads (threads associated with client connections). This table has a maximum size of 100 rows by default. To change the table size, modify the `performance_schema_setup_actors_size` system variable at server startup.

For each new foreground thread, the Performance Schema matches the user and host for the thread against the rows of the `setup_actors` table. If a row from that table matches, its `ENABLED` and `HISTORY` column values are used to set the the `INSTRUMENTED` and `HISTORY` columns, respectively, of the `threads` table row for the thread. This enables instrumenting and historical event logging to be applied selectively per host, user, or account (user and host combination). If there is no match, the `INSTRUMENTED` and `HISTORY` columns for the thread are set to `NO`.

For background threads, there is no associated user. `INSTRUMENTED` and `HISTORY` are `YES` by default and `setup_actors` is not consulted.

The initial contents of the `setup_actors` table match any user and host combination, so monitoring and historical event collection are enabled by default for all foreground threads:

```
mysql> SELECT * FROM setup_actors;
+------+------+------+---------+---------+
| HOST | USER | ROLE | ENABLED | HISTORY |
+------+------+------+---------+---------+
| %    | %    | %    | YES     | YES     |
+------+------+------+---------+---------+
```

For information about how to use the `setup_actors` table to affect event monitoring, see Section 25.4.6, “Pre-Filtering by Thread”.

Modifications to the `setup_actors` table affect only foreground threads created subsequent to the modification, not existing threads. To affect existing threads, modify the `INSTRUMENTED` and `HISTORY` columns of `threads` table rows.

The `setup_actors` table has these columns:

- `HOST`

  The host name. This should be a literal name, or `'%'` to mean “any host.”

- `USER`

  The user name. This should be a literal name, or `'%'` to mean “any user.”

- `ROLE`

  Unused.

- `ENABLED`

  Whether to enable instrumentation for foreground threads matched by the row. The value is `YES` or `NO`.

- `HISTORY`

  Whether to log historical events for foreground threads matched by the row. The value is `YES` or `NO`.

The `setup_actors` table has these indexes:

- Primary key on (`HOST`, `USER`, `ROLE`)

`TRUNCATE TABLE` is permitted for the `setup_actors` table. It removes the rows.

25.11.2.2 The setup_consumers Table

The `setup_consumers` table lists the types of consumers for which event information can be stored and which are enabled:

```
mysql> SELECT * FROM setup_consumers;
+--------------------------------+---------+
| NAME                           | ENABLED |
+--------------------------------+---------+
| events_stages_current          | NO      |
| events_stages_history          | NO      |
| events_stages_history_long     | NO      |
| events_statements_current      | YES     |
| events_statements_history      | YES     |
| events_statements_history_long | NO      |
| events_transactions_current    | YES     |
| events_transactions_history    | YES     |
| events_transactions_history_long | NO    |
| events_waits_current           | NO      |
| events_waits_history           | NO      |
| events_waits_history_long      | NO      |
| global_instrumentation         | YES     |
| thread_instrumentation         | YES     |
| statements_digest              | YES     |
+--------------------------------+---------+
```

The consumer settings in the `setup_consumers` table form a hierarchy from higher levels to lower. For detailed information about the effect of enabling different consumers, see Section 25.4.7, “Pre-Filtering by Consumer”.

Modifications to the `setup_consumers` table affect monitoring immediately.

The `setup_consumers` table has these columns:

- `NAME`

  The consumer name.

- `ENABLED`

  Whether the consumer is enabled. The value is `YES` or `NO`. This column can be modified. If you disable a consumer, the server does not spend time adding event information to it.

The `setup_consumers` table has these indexes:

- Primary key on (`NAME`)

`TRUNCATE TABLE` is not permitted for the `setup_consumers` table.

25.11.2.3 The setup_instruments Table

The `setup_instruments` table lists classes of instrumented objects for which events can be collected:

```
mysql> SELECT * FROM setup_instruments;
+-------------------------------------------------------------+---------+-------+
| NAME                                                        | ENABLED | TIMED |
+-------------------------------------------------------------+---------+-------+
...
| wait/synch/mutex/sql/LOCK_global_read_lock                  | YES     | YES   |
| wait/synch/mutex/sql/LOCK_global_system_variables           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_lock_db                           | YES     | YES   |
| wait/synch/mutex/sql/LOCK_manager                           | YES     | YES   |
...
| wait/synch/rwlock/sql/LOCK_grant                            | YES     | YES   |
| wait/synch/rwlock/sql/LOGGER::LOCK_logger                   | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_connect                 | YES     | YES   |
| wait/synch/rwlock/sql/LOCK_sys_init_slave                   | YES     | YES   |
...
| wait/io/file/sql/binlog                                     | YES     | YES   |
| wait/io/file/sql/binlog_index                               | YES     | YES   |
| wait/io/file/sql/casetest                                   | YES     | YES   |
| wait/io/file/sql/dbopt                                      | YES     | YES   |
...
```

Each instrument added to the source code provides a row for this table, even when the instrumented code is not executed. When an instrument is enabled and executed, instrumented instances are created, which are visible in the `*_instances` tables.

Modifications to most `setup_instruments` rows affect monitoring immediately. For some instruments, modifications are effective only at server startup; changing them at runtime has no effect. This affects primarily mutexes, conditions, and rwlocks in the server, although there may be other instruments for which this is true.

For more information about the role of the `setup_instruments` table in event filtering, see Section 25.4.3, “Event Pre-Filtering”.

The `setup_instruments` table has these columns:

- `NAME`

  The instrument name. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”. Events produced from execution of an instrument have an `EVENT_NAME` value that is taken from the instrument `NAME` value. (Events do not really have a “name,” but this provides a way to associate events with instruments.)

- `ENABLED`

  Whether the instrument is enabled. The value is `YES` or `NO`. This column can be modified. A disabled instrument produces no events.

- `TIMED`

  Whether the instrument is timed. This column can be modified.

  For memory instruments, the `TIMED` column in `setup_instruments` is ignored because memory operations are not timed.

  If an enabled instrument is not timed, the instrument code is enabled, but the timer is not. Events produced by the instrument have `NULL` for the `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` timer values. This in turn causes those values to be ignored when calculating the sum, minimum, maximum, and average time values in summary tables.

The `setup_instruments` table has these indexes:

- Primary key on (NAME)

TRUNCATE TABLE is not permitted for the setup_instruments table.

25.11.2.4 The setup_objects Table

The setup_objects table controls whether the Performance Schema monitors particular objects. This table has a maximum size of 100 rows by default. To change the table size, modify the performance_schema_setup_objects_size system variable at server startup.

The initial setup_objects contents look like this:

```
mysql> SELECT * FROM setup_objects;
+-------------+--------------------+-------------+---------+-------+
| OBJECT_TYPE | OBJECT_SCHEMA      | OBJECT_NAME | ENABLED | TIMED |
+-------------+--------------------+-------------+---------+-------+
| EVENT       | mysql              | %           | NO      | NO    |
| EVENT       | performance_schema | %           | NO      | NO    |
| EVENT       | information_schema | %           | NO      | NO    |
| EVENT       | %                  | %           | YES     | YES   |
| FUNCTION    | mysql              | %           | NO      | NO    |
| FUNCTION    | performance_schema | %           | NO      | NO    |
| FUNCTION    | information_schema | %           | NO      | NO    |
| FUNCTION    | %                  | %           | YES     | YES   |
| PROCEDURE   | mysql              | %           | NO      | NO    |
| PROCEDURE   | performance_schema | %           | NO      | NO    |
| PROCEDURE   | information_schema | %           | NO      | NO    |
| PROCEDURE   | %                  | %           | YES     | YES   |
| TABLE       | mysql              | %           | NO      | NO    |
| TABLE       | performance_schema | %           | NO      | NO    |
| TABLE       | information_schema | %           | NO      | NO    |
| TABLE       | %                  | %           | YES     | YES   |
| TRIGGER     | mysql              | %           | NO      | NO    |
| TRIGGER     | performance_schema | %           | NO      | NO    |
| TRIGGER     | information_schema | %           | NO      | NO    |
| TRIGGER     | %                  | %           | YES     | YES   |
+-------------+--------------------+-------------+---------+-------+
```

Modifications to the setup_objects table affect object monitoring immediately.

For object types listed in setup_objects, the Performance Schema uses the table to how to monitor them. Object matching is based on the OBJECT_SCHEMA and OBJECT_NAME columns. Objects for which there is no match are not monitored.

The effect of the default object configuration is to instrument all tables except those in the mysql, INFORMATION_SCHEMA, and performance_schema databases. (Tables in the INFORMATION_SCHEMA database are not instrumented regardless of the contents of setup_objects; the row for information_schema.% simply makes this default explicit.)

When the Performance Schema checks for a match in setup_objects, it tries to find more specific matches first. For example, with a table db1.t1, it looks for a match for 'db1' and 't1', then for 'db1' and '%', then for '%' and '%'. The order in which matching occurs matters because different matching setup_objects rows can have different ENABLED and TIMED values.

Rows can be inserted into or deleted from setup_objects by users with the INSERT or DELETE privilege on the table. For existing rows, only the ENABLED and TIMED columns can be modified, by users with the UPDATE privilege on the table.

For more information about the role of the setup_objects table in event filtering, see Section 25.4.3, “Event Pre-Filtering”.

The setup_objects table has these columns:

- OBJECT_TYPE

The type of object to instrument. The value is one of `'EVENT'` (Event Scheduler event), `'FUNCTION'` (stored function), `'PROCEDURE'` (stored procedure), `'TABLE'` (base table), or `'TRIGGER'` (trigger).

`TABLE` filtering affects table I/O events (`wait/io/table/sql/handler` instrument) and table lock events (`wait/lock/table/sql/handler` instrument).

- `OBJECT_SCHEMA`

  The schema that contains the object. This should be a literal name, or `'%'` to mean "any schema."

- `OBJECT_NAME`

  The name of the instrumented object. This should be a literal name, or `'%'` to mean "any object."

- `ENABLED`

  Whether events for the object are instrumented. The value is `YES` or `NO`. This column can be modified.

- `TIMED`

  Whether events for the object are timed. This column can be modified.

The `setup_objects` table has these indexes:

- Index on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`)

`TRUNCATE TABLE` is permitted for the `setup_objects` table. It removes the rows.

25.11.2.5 The setup_timers Table

The `setup_timers` table shows the currently selected event timers:

```
mysql> SELECT * FROM setup_timers;
+-------------+-------------+
| NAME        | TIMER_NAME  |
+-------------+-------------+
| idle        | MICROSECOND |
| wait        | CYCLE       |
| stage       | NANOSECOND  |
| statement   | NANOSECOND  |
| transaction | NANOSECOND  |
+-------------+-------------+
```

The `setup_timers.TIMER_NAME` value can be changed to select a different timer. The value can be any of the values in the `performance_timers.TIMER_NAME` column. For an explanation of how event timing occurs, see Section 25.4.1, "Performance Schema Event Timing".

Modifications to the `setup_timers` table affect monitoring immediately. Events already in progress may use the original timer for the begin time and the new timer for the end time. To avoid unpredictable results after you make timer changes, use `TRUNCATE TABLE` to reset Performance Schema statistics.

The `setup_timers` table has these columns:

- `NAME`

  The type of instrument the timer is used for.

- `TIMER_NAME`

  The timer that applies to the instrument type. This column can be modified.

The `setup_timers` table has these indexes:

- Primary key on (NAME)

TRUNCATE TABLE is not permitted for the setup_timers table.

25.11.3 Performance Schema Instance Tables

Instance tables document what types of objects are instrumented. They provide event names and explanatory notes or status information:

- cond_instances: Condition synchronization object instances
- file_instances: File instances
- mutex_instances: Mutex synchronization object instances
- rwlock_instances: Lock synchronization object instances
- socket_instances: Active connection instances

These tables list instrumented synchronization objects, files, and connections. There are three types of synchronization objects: cond, mutex, and rwlock. Each instance table has an EVENT_NAME or NAME column to indicate the instrument associated with each row. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

The mutex_instances.LOCKED_BY_THREAD_ID and rwlock_instances.WRITE_LOCKED_BY_THREAD_ID columns are extremely important for investigating performance bottlenecks or deadlocks. For examples of how to use them for this purpose, see Section 25.18, “Using the Performance Schema to Diagnose Problems”

25.11.3.1 The cond_instances Table

The cond_instances table lists all the conditions seen by the Performance Schema while the server executes. A condition is a synchronization mechanism used in the code to signal that a specific event has happened, so that a thread waiting for this condition can resume work.

When a thread is waiting for something to happen, the condition name is an indication of what the thread is waiting for, but there is no immediate way to tell which other thread, or threads, will cause the condition to happen.

The cond_instances table has these columns:

- NAME

  The instrument name associated with the condition.

- OBJECT_INSTANCE_BEGIN

  The address in memory of the instrumented condition.

The cond_instances table has these indexes:

- Primary key on (OBJECT_INSTANCE_BEGIN)
- Index on (NAME)

TRUNCATE TABLE is not permitted for the cond_instances table.

25.11.3.2 The file_instances Table

The file_instances table lists all the files seen by the Performance Schema when executing file I/O instrumentation. If a file on disk has never been opened, it will not be in file_instances. When a file is deleted from the disk, it is also removed from the file_instances table.

The `file_instances` table has these columns:

- `FILE_NAME`

  The file name.

- `EVENT_NAME`

  The instrument name associated with the file.

- `OPEN_COUNT`

  The count of open handles on the file. If a file was opened and then closed, it was opened 1 time, but `OPEN_COUNT` will be 0. To list all the files currently opened by the server, use `WHERE OPEN_COUNT > 0`.

The `file_instances` table has these indexes:

- Primary key on (`FILE_NAME`)
- Index on (`EVENT_NAME`)

`TRUNCATE TABLE` is not permitted for the `file_instances` table.

25.11.3.3 The mutex_instances Table

The `mutex_instances` table lists all the mutexes seen by the Performance Schema while the server executes. A mutex is a synchronization mechanism used in the code to enforce that only one thread at a given time can have access to some common resource. The resource is said to be “protected” by the mutex.

When two threads executing in the server (for example, two user sessions executing a query simultaneously) do need to access the same resource (a file, a buffer, or some piece of data), these two threads will compete against each other, so that the first query to obtain a lock on the mutex will cause the other query to wait until the first is done and unlocks the mutex.

The work performed while holding a mutex is said to be in a “critical section,” and multiple queries do execute this critical section in a serialized way (one at a time), which is a potential bottleneck.

The `mutex_instances` table has these columns:

- `NAME`

  The instrument name associated with the mutex.

- `OBJECT_INSTANCE_BEGIN`

  The address in memory of the instrumented mutex.

- `LOCKED_BY_THREAD_ID`

  When a thread currently has a mutex locked, `LOCKED_BY_THREAD_ID` is the `THREAD_ID` of the locking thread, otherwise it is `NULL`.

The `mutex_instances` table has these indexes:

- Primary key on (`OBJECT_INSTANCE_BEGIN`)
- Index on (`NAME`)
- Index on (`LOCKED_BY_THREAD_ID`)

`TRUNCATE TABLE` is not permitted for the `mutex_instances` table.

For every mutex instrumented in the code, the Performance Schema provides the following information.

- The `setup_instruments` table lists the name of the instrumentation point, with the prefix `wait/synch/mutex/`.
- When some code creates a mutex, a row is added to the `mutex_instances` table. The `OBJECT_INSTANCE_BEGIN` column is a property that uniquely identifies the mutex.
- When a thread attempts to lock a mutex, the `events_waits_current` table shows a row for that thread, indicating that it is waiting on a mutex (in the `EVENT_NAME` column), and indicating which mutex is waited on (in the `OBJECT_INSTANCE_BEGIN` column).
- When a thread succeeds in locking a mutex:
  - `events_waits_current` shows that the wait on the mutex is completed (in the `TIMER_END` and `TIMER_WAIT` columns)
  - The completed wait event is added to the `events_waits_history` and `events_waits_history_long` tables
  - `mutex_instances` shows that the mutex is now owned by the thread (in the `THREAD_ID` column).
- When a thread unlocks a mutex, `mutex_instances` shows that the mutex now has no owner (the `THREAD_ID` column is `NULL`).
- When a mutex object is destroyed, the corresponding row is removed from `mutex_instances`.

By performing queries on both of the following tables, a monitoring application or a DBA can detect bottlenecks or deadlocks between threads that involve mutexes:

- `events_waits_current`, to see what mutex a thread is waiting for
- `mutex_instances`, to see which other thread currently owns a mutex

25.11.3.4 The rwlock_instances Table

The `rwlock_instances` table lists all the rwlock (read write lock) instances seen by the Performance Schema while the server executes. An `rwlock` is a synchronization mechanism used in the code to enforce that threads at a given time can have access to some common resource following certain rules. The resource is said to be “protected” by the `rwlock`. The access is either shared (many threads can have a read lock at the same time), exclusive (only one thread can have a write lock at a given time), or shared-exclusive (a thread can have a write lock while permitting inconsistent reads by other threads). Shared-exclusive access is otherwise known as an `sxlock` and optimizes concurrency and improves scalability for read-write workloads.

Depending on how many threads are requesting a lock, and the nature of the locks requested, access can be either granted in shared mode, exclusive mode, shared-exclusive mode or not granted at all, waiting for other threads to finish first.

The `rwlock_instances` table has these columns:

- `NAME`

  The instrument name associated with the lock.

- `OBJECT_INSTANCE_BEGIN`

  The address in memory of the instrumented lock.

- `WRITE_LOCKED_BY_THREAD_ID`

When a thread currently has an `rwlock` locked in exclusive (write) mode, `WRITE_LOCKED_BY_THREAD_ID` is the `THREAD_ID` of the locking thread, otherwise it is `NULL`.

- `READ_LOCKED_BY_COUNT`

  When a thread currently has an `rwlock` locked in shared (read) mode, `READ_LOCKED_BY_COUNT` is incremented by 1. This is a counter only, so it cannot be used directly to find which thread holds a read lock, but it can be used to see whether there is a read contention on an `rwlock`, and see how many readers are currently active.

The `rwlock_instances` table has these indexes:

- Primary key on (`OBJECT_INSTANCE_BEGIN`)
- Index on (`NAME`)
- Index on (`WRITE_LOCKED_BY_THREAD_ID`)

`TRUNCATE TABLE` is not permitted for the `rwlock_instances` table.

By performing queries on both of the following tables, a monitoring application or a DBA may detect some bottlenecks or deadlocks between threads that involve locks:

- `events_waits_current`, to see what `rwlock` a thread is waiting for
- `rwlock_instances`, to see which other thread currently owns an `rwlock`

There is a limitation: The `rwlock_instances` can be used only to identify the thread holding a write lock, but not the threads holding a read lock.

25.11.3.5 The socket_instances Table

The `socket_instances` table provides a real-time snapshot of the active connections to the MySQL server. The table contains one row per TCP/IP or Unix socket file connection. Information available in this table provides a real-time snapshot of the active connections to the server. (Additional information is available in socket summary tables, including network activity such as socket operations and number of bytes transmitted and received; see Section 25.11.15.9, “Socket Summary Tables”).

```
mysql> SELECT * FROM socket_instances\G
*************************** 1. row ***************************
           EVENT_NAME: wait/io/socket/sql/server_unix_socket
OBJECT_INSTANCE_BEGIN: 4316619408
            THREAD_ID: 1
            SOCKET_ID: 16
                   IP:
                 PORT: 0
                STATE: ACTIVE
*************************** 2. row ***************************
           EVENT_NAME: wait/io/socket/sql/client_connection
OBJECT_INSTANCE_BEGIN: 4316644608
            THREAD_ID: 21
            SOCKET_ID: 39
                   IP: 127.0.0.1
                 PORT: 55233
                STATE: ACTIVE
*************************** 3. row ***************************
           EVENT_NAME: wait/io/socket/sql/server_tcpip_socket
OBJECT_INSTANCE_BEGIN: 4316699040
            THREAD_ID: 1
            SOCKET_ID: 14
                   IP: 0.0.0.0
                 PORT: 50603
                STATE: ACTIVE
```

Socket instruments have names of the form `wait/io/socket/sql/socket_type` and are used like this:

1. The server has a listening socket for each network protocol that it supports. The instruments associated with listening sockets for TCP/IP or Unix socket file connections have a *`socket_type`* value of `server_tcpip_socket` or `server_unix_socket`, respectively.
2. When a listening socket detects a connection, the server transfers the connection to a new socket managed by a separate thread. The instrument for the new connection thread has a *`socket_type`* value of `client_connection`.
3. When a connection terminates, the row in `socket_instances` corresponding to it is deleted.

The `socket_instances` table has these columns:

- `EVENT_NAME`

  The name of the `wait/io/socket/*` instrument that produced the event. This is a `NAME` value from the `setup_instruments` table. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

- `OBJECT_INSTANCE_BEGIN`

  This column uniquely identifies the socket. The value is the address of an object in memory.

- `THREAD_ID`

  The internal thread identifier assigned by the server. Each socket is managed by a single thread, so each socket can be mapped to a thread which can be mapped to a server process.

- `SOCKET_ID`

  The internal file handle assigned to the socket.

- `IP`

  The client IP address. The value may be either an IPv4 or IPv6 address, or blank to indicate a Unix socket file connection.

- `PORT`

  The TCP/IP port number, in the range from 0 to 65535.

- `STATE`

  The socket status, either `IDLE` or `ACTIVE`. Wait times for active sockets are tracked using the corresponding socket instrument. Wait times for idle sockets are tracked using the `idle` instrument.

  A socket is idle if it is waiting for a request from the client. When a socket becomes idle, the event row in `socket_instances` that is tracking the socket switches from a status of `ACTIVE` to `IDLE`. The `EVENT_NAME` value remains `wait/io/socket/*`, but timing for the instrument is suspended. Instead, an event is generated in the `events_waits_current` table with an `EVENT_NAME` value of `idle`.

  When the next request is received, the `idle` event terminates, the socket instance switches from `IDLE` to `ACTIVE`, and timing of the socket instrument resumes.

The `socket_instances` table has these indexes:

- Primary key on (`OBJECT_INSTANCE_BEGIN`)
- Index on (`THREAD_ID`)
- Index on (`SOCKET_ID`)

- Index on (`IP`, `PORT`)

`TRUNCATE TABLE` is not permitted for the `socket_instances` table.

The `IP:PORT` column combination value identifies the connection. This combination value is used in the `OBJECT_NAME` column of the `events_waits_xxx` tables, to identify the connection from which socket events come:

- For the Unix domain listener socket (`server_unix_socket`), the port is 0, and the IP is `''`.
- For client connections via the Unix domain listener (`client_connection`), the port is 0, and the IP is `''`.
- For the TCP/IP server listener socket (`server_tcpip_socket`), the port is always the master port (for example, 3306), and the IP is always `0.0.0.0`.
- For client connections via the TCP/IP listener (`client_connection`), the port is whatever the server assigns, but never 0. The IP is the IP of the originating host (`127.0.0.1` or `::1` for the local host)

25.11.4 Performance Schema Wait Event Tables

The Performance Schema instruments waits, which are events that take time. Within the event hierarchy, wait events nest within stage events, which nest within statement events, which nest within transaction events.

These tables store wait events:

- `events_waits_current`: Current wait events
- `events_waits_history`: The most recent wait events per thread
- `events_waits_history_long`: The most recent wait events globally (across all threads)

The following sections describe the wait event tables. There are also summary tables that aggregate information about wait events; see Section 25.11.15.1, “Wait Event Summary Tables”.

Wait Event Configuration

To control collection of wait events, set the state of the relevant instruments and consumers:

- The `setup_instruments` table contains instruments with names that begin with `wait`. Use these instruments to enable or disable collection of wait events.
- The `setup_consumers` table contains consumer values with names corresponding to the current and recent wait event table names. Use these consumers to filter collection of wait events.

Some wait instruments are enabled by default; others are disabled. For example:

```
mysql> SELECT * FROM setup_instruments
       WHERE NAME LIKE 'wait/io/file/innodb%';
+----------------------------------+---------+-------+
| NAME                             | ENABLED | TIMED |
+----------------------------------+---------+-------+
| wait/io/file/innodb/innodb_data_file | YES     | YES   |
| wait/io/file/innodb/innodb_log_file  | YES     | YES   |
| wait/io/file/innodb/innodb_temp_file | YES     | YES   |
+----------------------------------+---------+-------+
mysql> SELECT * FROM setup_instruments WHERE
       NAME LIKE 'wait/io/socket/%';
+---------------------------------------+---------+-------+
| NAME                                  | ENABLED | TIMED |
+---------------------------------------+---------+-------+
| wait/io/socket/sql/server_tcpip_socket | NO      | NO    |
| wait/io/socket/sql/server_unix_socket  | NO      | NO    |
```

```
| wait/io/socket/sql/client_connection   | NO      | NO    |
+----------------------------------------+---------+-------+
```

The wait consumers are disabled by default:

```
mysql> SELECT * FROM setup_consumers WHERE NAME LIKE '%waits%';
+--------------------------+---------+
| NAME                     | ENABLED |
+--------------------------+---------+
| events_waits_current     | NO      |
| events_waits_history     | NO      |
| events_waits_history_long | NO      |
+--------------------------+---------+
```

To control wait event collection at server startup, use lines like these in your `my.cnf` file:

- Enable:

```
[mysqld]
performance-schema-instrument='wait/%=ON'
performance-schema-consumer-events-waits-current=ON
performance-schema-consumer-events-waits-history=ON
performance-schema-consumer-events-waits-history-long=ON
```

- Disable:

```
[mysqld]
performance-schema-instrument='wait/%=OFF'
performance-schema-consumer-events-waits-current=OFF
performance-schema-consumer-events-waits-history=OFF
performance-schema-consumer-events-waits-history-long=OFF
```

To control wait event collection at runtime, update the `setup_instruments` and `setup_consumers` tables:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
WHERE NAME = 'wait/%';
UPDATE setup_consumers SET ENABLED = 'YES'
WHERE NAME LIKE '%waits%';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO', TIMED = 'NO'
WHERE NAME = 'wait/%';
UPDATE setup_consumers SET ENABLED = 'NO'
WHERE NAME LIKE '%waits%';
```

To collect only specific wait events, enable only the corresponding wait instruments. To collect wait events only for specific wait event tables, enable the wait instruments but only the wait consumers corresponding to the desired tables.

The `setup_timers` table contains a row with a `NAME` value of `wait` that indicates the unit for wait event timing. The default unit is `CYCLE`:

```
mysql> SELECT * FROM setup_timers WHERE NAME = 'wait';
+------+------------+
| NAME | TIMER_NAME |
+------+------------+
| wait | CYCLE      |
+------+------------+
```

To change the timing unit, modify the `TIMER_NAME` value:

```
UPDATE setup_timers SET TIMER_NAME = 'NANOSECOND'
WHERE NAME = 'wait';
```

For additional information about configuring event collection, see Section 25.3, “Performance Schema Startup Configuration”, and Section 25.4, “Performance Schema Runtime Configuration”.

25.11.4.1 The events_waits_current Table

The `events_waits_current` table contains current wait events, one row per thread showing the current status of the thread's most recent monitored wait event.

Of the tables that contain wait event rows, `events_waits_current` is the most fundamental. Other tables that contain wait event rows are logically derived from the current events. For example, the `events_waits_history` and `events_waits_history_long` tables are collections of the most recent wait events, up to a fixed number of rows.

For information about configuration of wait event collection, see Section 25.11.4, “Performance Schema Wait Event Tables”.

The `events_waits_current` table has these columns:

- `THREAD_ID`, `EVENT_ID`

  The thread associated with the event and the thread current event number when the event starts. The `THREAD_ID` and `EVENT_ID` values taken together uniquely identify the row. No two rows have the same pair of values.

- `END_EVENT_ID`

  This column is set to `NULL` when the event starts and updated to the thread current event number when the event ends.

- `EVENT_NAME`

  The name of the instrument that produced the event. This is a `NAME` value from the `setup_instruments` table. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

- `SOURCE`

  The name of the source file containing the instrumented code that produced the event and the line number in the file at which the instrumentation occurs. This enables you to check the source to determine exactly what code is involved. For example, if a mutex or lock is being blocked, you can check the context in which this occurs.

- `TIMER_START`, `TIMER_END`, `TIMER_WAIT`

  Timing information for the event. The unit for these values is picoseconds (trillionths of a second). The `TIMER_START` and `TIMER_END` values indicate when event timing started and ended. `TIMER_WAIT` is the event elapsed time (duration).

  If an event has not finished, `TIMER_END` is the current timer value and `TIMER_WAIT` is the time elapsed so far (`TIMER_END` − `TIMER_START`).

  If an event is produced from an instrument that has `TIMED = NO`, timing information is not collected, and `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` are all `NULL`.

  For discussion of picoseconds as the unit for event times and factors that affect time values, see Section 25.4.1, “Performance Schema Event Timing”.

- `SPINS`

For a mutex, the number of spin rounds. If the value is `NULL`, the code does not use spin rounds or spinning is not instrumented.

- `OBJECT_SCHEMA`, `OBJECT_NAME`, `OBJECT_TYPE`, `OBJECT_INSTANCE_BEGIN`

  These columns identify the object "being acted on." What that means depends on the object type.

  For a synchronization object (`cond`, `mutex`, `rwlock`):

  - `OBJECT_SCHEMA`, `OBJECT_NAME`, and `OBJECT_TYPE` are `NULL`.
  - `OBJECT_INSTANCE_BEGIN` is the address of the synchronization object in memory.

  For a file I/O object:

  - `OBJECT_SCHEMA` is `NULL`.
  - `OBJECT_NAME` is the file name.
  - `OBJECT_TYPE` is `FILE`.
  - `OBJECT_INSTANCE_BEGIN` is an address in memory.

  For a socket object:

  - `OBJECT_NAME` is the `IP:PORT` value for the socket.
  - `OBJECT_INSTANCE_BEGIN` is an address in memory.

  For a table I/O object:

  - `OBJECT_SCHEMA` is the name of the schema that contains the table.
  - `OBJECT_NAME` is the table name.
  - `OBJECT_TYPE` is `TABLE` for a persistent base table or `TEMPORARY TABLE` for a temporary table.
  - `OBJECT_INSTANCE_BEGIN` is an address in memory.

  An `OBJECT_INSTANCE_BEGIN` value itself has no meaning, except that different values indicate different objects. `OBJECT_INSTANCE_BEGIN` can be used for debugging. For example, it can be used with `GROUP BY OBJECT_INSTANCE_BEGIN` to see whether the load on 1,000 mutexes (that protect, say, 1,000 pages or blocks of data) is spread evenly or just hitting a few bottlenecks. This can help you correlate with other sources of information if you see the same object address in a log file or another debugging or performance tool.

- `INDEX_NAME`

  The name of the index used. `PRIMARY` indicates the table primary index. `NULL` means that no index was used.

- `NESTING_EVENT_ID`

  The `EVENT_ID` value of the event within which this event is nested.

- `NESTING_EVENT_TYPE`

  The nesting event type. The value is `TRANSACTION`, `STATEMENT`, `STAGE`, or `WAIT`.

- `OPERATION`

  The type of operation performed, such as `lock`, `read`, or `write`.

- NUMBER_OF_BYTES

  The number of bytes read or written by the operation. For table I/O waits (events for the wait/io/table/sql/handler instrument), NUMBER_OF_BYTES indicates the number of rows. If the value is greater than 1, the event is for a batch I/O operation. The following discussion describes the difference between exclusively single-row reporting and reporting that reflects batch I/O.

  MySQL executes joins using a nested-loop implementation. The job of the Performance Schema instrumentation is to provide row count and accumulated execution time per table in the join. Assume a join query of the following form that is executed using a table join order of t1, t2, t3:

  ```
  SELECT ... FROM t1 JOIN t2 ON ... JOIN t3 ON ...
  ```

  Table “fanout” is the increase or decrease in number of rows from adding a table during join processing. If the fanout for table t3 is greater than 1, the majority of row-fetch operations are for that table. Suppose that the join accesses 10 rows from t1, 20 rows from t2 per row from t1, and 30 rows from t3 per row of table t2. With single-row reporting, the total number of instrumented operations is:

  ```
  10 + (10 * 20) + (10 * 20 * 30) = 6210
  ```

  A significant reduction in the number of instrumented operations is achievable by aggregating them per scan (that is, per unique combination of rows from t1 and t2). With batch I/O reporting, the Performance Schema produces an event for each scan of the innermost table t3 rather than for each row, and the number of instrumented row operations reduces to:

  ```
  10 + (10 * 20) + (10 * 20) = 410
  ```

  That is a reduction of 93%, illustrating how the batch-reporting strategy significantly reduces Performance Schema overhead for table I/O by reducing the number of reporting calls. The tradeoff is lesser accuracy for event timing. Rather than time for an individual row operation as in per-row reporting, timing for batch I/O includes time spent for operations such as join buffering, aggregation, and returning rows to the client.

  For batch I/O reporting to occur, these conditions must be true:

  - Query execution accesses the innermost table of a query block (for a single-table query, that table counts as innermost)
  - Query execution does not request a single row from the table (so, for example, eq_ref access prevents use of batch reporting)
  - Query execution does not evaluate a subquery containing table access for the table

- FLAGS

  Reserved for future use.

The events_waits_current table has these indexes:

- Primary key on (THREAD_ID, EVENT_ID)

TRUNCATE TABLE is permitted for the events_waits_current table. It removes the rows.

25.11.4.2 The events_waits_history Table

The events_waits_history table contains the most recent *N* wait events per thread. The value of *N* is autosized at server startup. To set the table size explicitly, set the performance_schema_events_waits_history_size system variable at server startup. Wait

events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full.

The `events_waits_history` table has the same structure as `events_waits_current`, including indexing. See Section 25.11.4.1, “The events_waits_current Table”.

`TRUNCATE TABLE` is permitted for the `events_waits_history` table. It removes the rows.

For information about configuration of wait event collection, see Section 25.11.4, “Performance Schema Wait Event Tables”.

25.11.4.3 The events_waits_history_long Table

The `events_waits_history_long` table contains the most recent *`N`* wait events. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_waits_history_long_size` system variable at server startup. Wait events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full. When a thread ends, its rows are removed from the table.

The `events_waits_history_long` table has the same structure as `events_waits_current`, except that it has no indexes. See Section 25.11.4.1, “The events_waits_current Table”.

`TRUNCATE TABLE` is permitted for the `events_waits_history_long` table. It removes the rows.

For information about configuration of wait event collection, see Section 25.11.4, “Performance Schema Wait Event Tables”.

25.11.5 Performance Schema Stage Event Tables

The Performance Schema instruments stages, which are steps during the statement-execution process, such as parsing a statement, opening a table, or performing a `filesort` operation. Stages correspond to the thread states displayed by `SHOW PROCESSLIST` or that are visible in the `INFORMATION_SCHEMA.PROCESSLIST` table. Stages begin and end when state values change.

Within the event hierarchy, wait events nest within stage events, which nest within statement events, which nest within transaction events.

These tables store stage events:

- `events_stages_current`: Current stage events
- `events_stages_history`: The most recent stage events per thread
- `events_stages_history_long`: The most recent stage events globally (across all threads)

The following sections describe the stage event tables. There are also summary tables that aggregate information about stage events; see Section 25.11.15.2, “Stage Summary Tables”.

Stage Event Configuration

To control collection of stage events, set the state of the relevant instruments and consumers:

- The `setup_instruments` table contains instruments with names that begin with `stage`. Use these instruments to enable or disable collection of stage events.
- The `setup_consumers` table contains consumer values with names corresponding to the current and recent stage event table names. Use these consumers to filter collection of stage events.

Other than those instruments that provide statement progress information, the stage instruments are disabled by default. For example:

```
mysql> SELECT * FROM setup_instruments WHERE NAME RLIKE 'stage/sql/[a-c]';
+-----------------------------------------------------+---------+-------+
| NAME                                                | ENABLED | TIMED |
+-----------------------------------------------------+---------+-------+
| stage/sql/After create                              | NO      | NO    |
| stage/sql/allocating local table                    | NO      | NO    |
| stage/sql/altering table                            | NO      | NO    |
| stage/sql/committing alter table to storage engine | NO      | NO    |
| stage/sql/Changing master                           | NO      | NO    |
| stage/sql/Checking master version                   | NO      | NO    |
| stage/sql/checking permissions                      | NO      | NO    |
| stage/sql/checking privileges on cached query       | NO      | NO    |
| stage/sql/checking query cache for query            | NO      | NO    |
| stage/sql/cleaning up                               | NO      | NO    |
| stage/sql/closing tables                            | NO      | NO    |
| stage/sql/Connecting to master                      | NO      | NO    |
| stage/sql/converting HEAP to MyISAM                 | NO      | NO    |
| stage/sql/Copying to group table                    | NO      | NO    |
| stage/sql/Copying to tmp table                      | NO      | NO    |
| stage/sql/copy to tmp table                         | NO      | NO    |
| stage/sql/Creating sort index                       | NO      | NO    |
| stage/sql/creating table                            | NO      | NO    |
| stage/sql/Creating tmp table                        | NO      | NO    |
+-----------------------------------------------------+---------+-------+
```

Stage event instruments that provide statement progress information are enabled and timed by default:

```
mysql> SELECT * FROM setup_instruments WHERE
       ENABLED='YES' AND NAME LIKE "stage/%";
+-----------------------------------------------------+---------+-------+
| NAME                                                | ENABLED | TIMED |
+-----------------------------------------------------+---------+-------+
| stage/sql/copy to tmp table                         | YES     | YES   |
| stage/innodb/alter table (end)                      | YES     | YES   |
| stage/innodb/alter table (flush)                    | YES     | YES   |
| stage/innodb/alter table (insert)                   | YES     | YES   |
| stage/innodb/alter table (log apply index)          | YES     | YES   |
| stage/innodb/alter table (log apply table)          | YES     | YES   |
| stage/innodb/alter table (merge sort)               | YES     | YES   |
| stage/innodb/alter table (read PK and internal sort) | YES     | YES   |
| stage/innodb/buffer pool load                       | YES     | YES   |
+-----------------------------------------------------+---------+-------+
```

The stage consumers are disabled by default:

```
mysql> SELECT * FROM setup_consumers WHERE NAME LIKE '%stages%';
+----------------------------+---------+
| NAME                       | ENABLED |
+----------------------------+---------+
| events_stages_current      | NO      |
| events_stages_history      | NO      |
| events_stages_history_long | NO      |
+----------------------------+---------+
```

To control stage event collection at server startup, use lines like these in your `my.cnf` file:

- Enable:

  ```
  [mysqld]
  performance-schema-instrument='stage/%=ON'
  performance-schema-consumer-events-stages-current=ON
  performance-schema-consumer-events-stages-history=ON
  performance-schema-consumer-events-stages-history-long=ON
  ```

- Disable:

```
[mysqld]
performance-schema-instrument='stage/%=OFF'
performance-schema-consumer-events-stages-current=OFF
performance-schema-consumer-events-stages-history=OFF
performance-schema-consumer-events-stages-history-long=OFF
```

To control stage event collection at runtime, update the `setup_instruments` and `setup_consumers` tables:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
WHERE NAME = 'stage/%';
UPDATE setup_consumers SET ENABLED = 'YES'
WHERE NAME LIKE '%stages%';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO', TIMED = 'NO'
WHERE NAME = 'stage/%';
UPDATE setup_consumers SET ENABLED = 'NO'
WHERE NAME LIKE '%stages%';
```

To collect only specific stage events, enable only the corresponding stage instruments. To collect stage events only for specific stage event tables, enable the stage instruments but only the stage consumers corresponding to the desired tables.

The `setup_timers` table contains a row with a `NAME` value of `stage` that indicates the unit for stage event timing. The default unit is `NANOSECOND`:

```
mysql> SELECT * FROM setup_timers WHERE NAME = 'stage';
+-------+------------+
| NAME  | TIMER_NAME |
+-------+------------+
| stage | NANOSECOND |
+-------+------------+
```

To change the timing unit, modify the `TIMER_NAME` value:

```
UPDATE setup_timers SET TIMER_NAME = 'MICROSECOND'
WHERE NAME = 'stage';
```

For additional information about configuring event collection, see Section 25.3, “Performance Schema Startup Configuration”, and Section 25.4, “Performance Schema Runtime Configuration”.

Stage Event Progress Information

The Performance Schema stage event tables contain two columns that, taken together, provide a stage progress indicator for each row:

- `WORK_COMPLETED`: The number of work units completed for the stage
- `WORK_ESTIMATED`: The number of work units expected for the stage

Each column is `NULL` if no progress information is provided for an instrument. Interpretation of the information, if it is available, depends entirely on the instrument implementation. The Performance Schema tables provide a container to store progress data, but make no assumptions about the semantics of the metric itself:

- A “work unit” is an integer metric that increases over time during execution, such as the number of bytes, rows, files, or tables processed. The definition of “work unit” for a particular instrument is left to the instrumentation code providing the data.

- The `WORK_COMPLETED` value can increase one or many units at a time, depending on the instrumented code.
- The `WORK_ESTIMATED` value can change during the stage, depending on the instrumented code.

Instrumentation for a stage event progress indicator can implement any of the following behaviors:

- No progress instrumentation

  This is the most typical case, where no progress data is provided. The `WORK_COMPLETED` and `WORK_ESTIMATED` columns are both `NULL`.

- Unbounded progress instrumentation

  Only the `WORK_COMPLETED` column is meaningful. No data is provided for the `WORK_ESTIMATED` column, which displays 0.

  By querying the `events_stages_current` table for the monitored session, a monitoring application can report how much work has been performed so far, but cannot report whether the stage is near completion. Currently, no stages are instrumented like this.

- Bounded progress instrumentation

  The `WORK_COMPLETED` and `WORK_ESTIMATED` columns are both meaningful.

  This type of progress indicator is appropriate for an an operation with a defined completion criterion, such as the table-copy instrument described later. By querying the `events_stages_current` table for the monitored session, a monitoring application can report how much work has been performed so far, and can report the overall completion percentage for the stage, by computing the `WORK_COMPLETED` / `WORK_ESTIMATED` ratio.

The `stage/sql/copy to tmp table` instrument illustrates how progress indicators work. During execution of an `ALTER TABLE` statement, the `stage/sql/copy to tmp table` stage is used, and this stage can execute potentially for a long time, depending on the size of the data to copy.

The table-copy task has a defined termination (all rows copied), and the `stage/sql/copy to tmp table` stage is instrumented to provided bounded progress information: The work unit used is number of rows copied, `WORK_COMPLETED` and `WORK_ESTIMATED` are both meaningful, and their ratio indicates task percentage complete.

To enable the instrument and the relevant consumers, execute these statements:

```
UPDATE setup_instruments SET ENABLED='YES'
WHERE NAME='stage/sql/copy to tmp table';
UPDATE setup_consumers SET ENABLED='YES'
WHERE NAME LIKE 'events_stages_%';
```

To see the progress of an ongoing `ALTER TABLE` statement, select from the `events_stages_current` table.

25.11.5.1 The events_stages_current Table

The `events_stages_current` table contains current stage events, one row per thread showing the current status of the thread's most recent monitored stage event.

Of the tables that contain stage event rows, `events_stages_current` is the most fundamental. Other tables that contain stage event rows are logically derived from the current events. For example, the `events_stages_history` and `events_stages_history_long` tables are collections of the most recent stage events, up to a fixed number of rows.

For information about configuration of stage event collection, see Section 25.11.5, “Performance Schema Stage Event Tables”.

The `events_stages_current` table has these columns:

- `THREAD_ID`, `EVENT_ID`

  The thread associated with the event and the thread current event number when the event starts. The `THREAD_ID` and `EVENT_ID` values taken together uniquely identify the row. No two rows have the same pair of values.

- `END_EVENT_ID`

  This column is set to `NULL` when the event starts and updated to the thread current event number when the event ends.

- `EVENT_NAME`

  The name of the instrument that produced the event. This is a `NAME` value from the `setup_instruments` table. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

- `SOURCE`

  The name of the source file containing the instrumented code that produced the event and the line number in the file at which the instrumentation occurs. This enables you to check the source to determine exactly what code is involved.

- `TIMER_START`, `TIMER_END`, `TIMER_WAIT`

  Timing information for the event. The unit for these values is picoseconds (trillionths of a second). The `TIMER_START` and `TIMER_END` values indicate when event timing started and ended. `TIMER_WAIT` is the event elapsed time (duration).

  If an event has not finished, `TIMER_END` is the current timer value and `TIMER_WAIT` is the time elapsed so far (`TIMER_END` − `TIMER_START`).

  If an event is produced from an instrument that has `TIMED = NO`, timing information is not collected, and `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` are all `NULL`.

  For discussion of picoseconds as the unit for event times and factors that affect time values, see Section 25.4.1, “Performance Schema Event Timing”.

- `WORK_COMPLETED`, `WORK_ESTIMATED`

  These columns provide stage progress information, for instruments that have been implemented to produce such information. `WORK_COMPLETED` indicates how many work units have been completed for the stage, and `WORK_ESTIMATED` indicates how many work units are expected for the stage. For more information, see Stage Event Progress Information.

- `NESTING_EVENT_ID`

  The `EVENT_ID` value of the event within which this event is nested. The nesting event for a stage event is usually a statement event.

- `NESTING_EVENT_TYPE`

  The nesting event type. The value is `TRANSACTION`, `STATEMENT`, `STAGE`, or `WAIT`.

The `events_stages_current` table has these indexes:

- Primary key on (`THREAD_ID`, `EVENT_ID`)

`TRUNCATE TABLE` is permitted for the `events_stages_current` table. It removes the rows.

25.11.5.2 The events_stages_history Table

The `events_stages_history` table contains the most recent *`N`* stage events per thread. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_stages_history_size` system variable at server startup. Stage events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full.

The `events_stages_history` table has the same structure as `events_stages_current`, including indexing. See Section 25.11.5.1, "The events_stages_current Table".

`TRUNCATE TABLE` is permitted for the `events_stages_history` table. It removes the rows.

For information about configuration of stage event collection, see Section 25.11.5, "Performance Schema Stage Event Tables".

25.11.5.3 The events_stages_history_long Table

The `events_stages_history_long` table contains the most recent *`N`* stage events. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_stages_history_long_size` system variable at server startup. Stage events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full. When a thread ends, its rows are removed from the table.

The `events_stages_history_long` table has the same structure as `events_stages_current`, except that it has no indexes. See Section 25.11.5.1, "The events_stages_current Table".

`TRUNCATE TABLE` is permitted for the `events_stages_history_long` table. It removes the rows.

For information about configuration of stage event collection, see Section 25.11.5, "Performance Schema Stage Event Tables".

25.11.6 Performance Schema Statement Event Tables

The Performance Schema instruments statement execution. Statement events occur at a high level of the event hierarchy. Within the event hierarchy, wait events nest within stage events, which nest within statement events, which nest within transaction events.

These tables store statement events:

- `events_statements_current`: Current statement events
- `events_statements_history`: The most recent statement events per thread
- `events_statements_history_long`: The most recent statement events globally (across all threads)
- `prepared_statements_instances`: Prepared statement instances and statistics

The following sections describe the statement event tables. There are also summary tables that aggregate information about statement events; see Section 25.11.15.3, "Statement Summary Tables".

Statement Event Configuration

To control collection of statement events, set the state of the relevant instruments and consumers:

- The `setup_instruments` table contains instruments with names that begin with `statement`. Use these instruments to enable or disable collection of statement events.
- The `setup_consumers` table contains consumer values with names corresponding to the current and recent statement event table names, and the statement digest consumer. Use these consumers to filter collection of statement events and statement digesting.

The statement instruments are enabled by default, and the `events_statements_current`, `events_statements_history`, and `statements_digest` statement consumers are enabled by default:

```
mysql> SELECT * FROM setup_instruments WHERE NAME LIKE 'statement/%';
+----------------------------------------+---------+-------+
| NAME                                   | ENABLED | TIMED |
+----------------------------------------+---------+-------+
| statement/sql/select                   | YES     | YES   |
| statement/sql/create_table             | YES     | YES   |
| statement/sql/create_index             | YES     | YES   |
...
| statement/sp/stmt                      | YES     | YES   |
| statement/sp/set                       | YES     | YES   |
| statement/sp/set_trigger_field         | YES     | YES   |
| statement/scheduler/event              | YES     | YES   |
| statement/com/Sleep                    | YES     | YES   |
| statement/com/Quit                     | YES     | YES   |
| statement/com/Init DB                  | YES     | YES   |
...
| statement/abstract/Query               | YES     | YES   |
| statement/abstract/new_packet          | YES     | YES   |
| statement/abstract/relay_log           | YES     | YES   |
+----------------------------------------+---------+-------+
```

```
mysql> SELECT * FROM setup_consumers WHERE NAME LIKE '%statements%';
+--------------------------------+---------+
| NAME                           | ENABLED |
+--------------------------------+---------+
| events_statements_current      | YES     |
| events_statements_history      | YES     |
| events_statements_history_long | NO      |
| statements_digest              | YES     |
+--------------------------------+---------+
```

To control statement event collection at server startup, use lines like these in your `my.cnf` file:

- Enable:

```
[mysqld]
performance-schema-instrument='statement/%=ON'
performance-schema-consumer-events-statements-current=ON
performance-schema-consumer-events-statements-history=ON
performance-schema-consumer-events-statements-history-long=ON
performance-schema-consumer-statements-digest=ON
```

- Disable:

```
[mysqld]
performance-schema-instrument='statement/%=OFF'
performance-schema-consumer-events-statements-current=OFF
performance-schema-consumer-events-statements-history=OFF
performance-schema-consumer-events-statements-history-long=OFF
performance-schema-consumer-statements-digest=OFF
```

To control statement event collection at runtime, update the `setup_instruments` and `setup_consumers` tables:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
WHERE NAME LIKE 'statement/%';
UPDATE setup_consumers SET ENABLED = 'YES'
WHERE NAME LIKE '%statements%';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO', TIMED = 'NO'
WHERE NAME LIKE 'statement/%';
UPDATE setup_consumers SET ENABLED = 'NO'
WHERE NAME LIKE '%statements%';
```

To collect only specific statement events, enable only the corresponding statement instruments. To collect statement events only for specific statement event tables, enable the statement instruments but only the statement consumers corresponding to the desired tables.

The `setup_timers` table contains a row with a `NAME` value of `statement` that indicates the unit for statement event timing. The default unit is `NANOSECOND`:

```
mysql> SELECT * FROM setup_timers WHERE NAME = 'statement';
+-----------+------------+
| NAME      | TIMER_NAME |
+-----------+------------+
| statement | NANOSECOND |
+-----------+------------+
```

To change the timing unit, modify the `TIMER_NAME` value:

```
UPDATE setup_timers SET TIMER_NAME = 'MICROSECOND'
WHERE NAME = 'statement';
```

For additional information about configuring event collection, see Section 25.3, “Performance Schema Startup Configuration”, and Section 25.4, “Performance Schema Runtime Configuration”.

Statement Monitoring

Statement monitoring begins from the moment the server sees that activity is requested on a thread, to the moment when all activity has ceased. Typically, this means from the time the server gets the first packet from the client to the time the server has finished sending the response. Statements within stored programs are monitored like other statements.

When the Performance Schema instruments a request (server command or SQL statement), it uses instrument names that proceed in stages from more general (or “abstract”) to more specific until it arrives at a final instrument name.

Final instrument names correspond to server commands and SQL statements:

- Server commands correspond to the `COM_xxx codes` defined in the `mysql_com.h` header file and processed in `sql/sql_parse.cc`. Examples are `COM_PING` and `COM_QUIT`. Instruments for commands have names that begin with `statement/com`, such as `statement/com/Ping` and `statement/com/Quit`.
- SQL statements are expressed as text, such as `DELETE FROM t1` or `SELECT * FROM t2`. Instruments for SQL statements have names that begin with `statement/sql`, such as `statement/sql/delete` and `statement/sql/select`.

Some final instrument names are specific to error handling:

- `statement/com/Error` accounts for messages received by the server that are out of band. It can be used to detect commands sent by clients that the server does not understand. This may be helpful for purposes such as identifying clients that are misconfigured or using a version of MySQL more recent than that of the server, or clients that are attempting to attack the server.
- `statement/sql/error` accounts for SQL statements that fail to parse. It can be used to detect malformed queries sent by clients. A query that fails to parse differs from a query that parses but fails due to an error during execution. For example, `SELECT * FROM` is malformed, and the `statement/sql/error` instrument is used. By contrast, `SELECT *` parses but fails with a `No`

`tables used` error. In this case, `statement/sql/select` is used and the statement event contains information to indicate the nature of the error.

A request can be obtained from any of these sources:

- As a command or statement request from a client, which sends the request as packets
- As a statement string read from the relay log on a replication slave
- As an event from the Event Scheduler

The details for a request are not initially known and the Performance Schema proceeds from abstract to specific instrument names in a sequence that depends on the source of the request.

For a request received from a client:

1. When the server detects a new packet at the socket level, a new statement is started with an abstract instrument name of `statement/abstract/new_packet`.
2. When the server reads the packet number, it knows more about the type of request received, and the Performance Schema refines the instrument name. For example, if the request is a `COM_PING` packet, the instrument name becomes `statement/com/Ping` and that is the final name. If the request is a `COM_QUERY` packet, it is known to correspond to an SQL statement but not the particular type of statement. In this case, the instrument changes from one abstract name to a more specific but still abstract name, `statement/abstract/Query`, and the request requires further classification.
3. If the request is a statement, the statement text is read and given to the parser. After parsing, the exact statement type is known. If the request is, for example, an `INSERT` statement, the Performance Schema refines the instrument name from `statement/abstract/Query` to `statement/sql/insert`, which is the final name.

For a request read as a statement from the relay log on a replication slave:

1. Statements in the relay log are stored as text and are read as such. There is no network protocol, so the `statement/abstract/new_packet` instrument is not used. Instead, the initial instrument is `statement/abstract/relay_log`.
2. When the statement is parsed, the exact statement type is known. If the request is, for example, an `INSERT` statement, the Performance Schema refines the instrument name from `statement/abstract/Query` to `statement/sql/insert`, which is the final name.

The preceding description applies only for statement-based replication. For row-based replication, table I/O done on the slave as it processes row changes can be instrumented, but row events in the relay log do not appear as discrete statements.

For a request received from the Event Scheduler:

The event execution is instrumented using the name `statement/scheduler/event`. This is the final name.

Statements executed within the event body are instrumented using `statement/sql/*` names, without use of any preceding abstract instrument. An event is a stored program, and stored programs are precompiled in memory before execution. Consequently, there is no parsing at runtime and the type of each statement is known by the time it executes.

Statements executed within the event body are child statements. For example, if an event executes an `INSERT` statement, execution of the event itself is the parent, instrumented using `statement/scheduler/event`, and the `INSERT` is the child, instrumented using `statement/sql/insert`. The parent/child relationship holds *between* separate instrumented operations. This differs from the sequence of refinement that occurs *within* a single instrumented operation, from abstract to final instrument names.

For statistics to be collected for statements, it is not sufficient to enable only the final `statement/sql/*` instruments used for individual statement types. The abtract `statement/abstract/*` instruments must be enabled as well. This should not normally be an issue because all statement instruments are enabled by default. However, an application that enables or disables statement instruments selectively must take into account that disabling abstract instruments also disables statistics collection for the individual statement instruments. For example, to collect statistics for `INSERT` statements, `statement/sql/insert` must be enabled, but also `statement/abstract/new_packet` and `statement/abstract/Query`. Similarly, for replicated statements to be instrumented, `statement/abstract/relay_log` must be enabled.

No statistics are aggregated for abstract instruments such as `statement/abstract/Query` because no statement is ever classified with an abstract instrument as the final statement name.

25.11.6.1 The events_statements_current Table

The `events_statements_current` table contains current statement events, one row per thread showing the current status of the thread's most recent monitored statement event.

Of the tables that contain statement event rows, `events_statements_current` is the most fundamental. Other tables that contain statement event rows are logically derived from the current events. For example, the `events_statements_history` and `events_statements_history_long` tables are collections of the most recent statement events, up to a fixed number of rows.

For information about configuration of statement event collection, see Section 25.11.6, “Performance Schema Statement Event Tables”.

The `events_statements_current` table has these columns:

- `THREAD_ID`, `EVENT_ID`

  The thread associated with the event and the thread current event number when the event starts. The `THREAD_ID` and `EVENT_ID` values taken together uniquely identify the row. No two rows have the same pair of values.

- `END_EVENT_ID`

  This column is set to `NULL` when the event starts and updated to the thread current event number when the event ends.

- `EVENT_NAME`

  The name of the instrument from which the event was collected. This is a `NAME` value from the `setup_instruments` table. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

  For SQL statements, the `EVENT_NAME` value initially is `statement/com/Query` until the statement is parsed, then changes to a more appropriate value, as described in Section 25.11.6, “Performance Schema Statement Event Tables”.

- `SOURCE`

  The name of the source file containing the instrumented code that produced the event and the line number in the file at which the instrumentation occurs. This enables you to check the source to determine exactly what code is involved.

- `TIMER_START`, `TIMER_END`, `TIMER_WAIT`

  Timing information for the event. The unit for these values is picoseconds (trillionths of a second). The `TIMER_START` and `TIMER_END` values indicate when event timing started and ended. `TIMER_WAIT` is the event elapsed time (duration).

If an event has not finished, `TIMER_END` is the current timer value and `TIMER_WAIT` is the time elapsed so far (`TIMER_END` − `TIMER_START`).

If an event is produced from an instrument that has `TIMED = NO`, timing information is not collected, and `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` are all `NULL`.

For discussion of picoseconds as the unit for event times and factors that affect time values, see Section 25.4.1, “Performance Schema Event Timing”.

- `LOCK_TIME`

  The time spent waiting for table locks. This value is computed in microseconds but normalized to picoseconds for easier comparison with other Performance Schema timers.

- `SQL_TEXT`

  The text of the SQL statement. For a command not associated with an SQL statement, the value is `NULL`.

  The maximum space available for statement display is 1024 bytes by default. To change this value, set the `performance_schema_max_sql_text_length` system variable at server startup.

- `DIGEST`

  The statement digest MD5 value as a string of 32 hexadecimal characters, or `NULL` if the `statements_digest` consumer is `no`. For more information about statement digesting, see Section 25.9, “Performance Schema Statement Digests”.

- `DIGEST_TEXT`

  The normalized statement digest text, or `NULL` if the `statements_digest` consumer is `no`. For more information about statement digesting, see Section 25.9, “Performance Schema Statement Digests”.

  The `performance_schema_max_digest_length` system variable determines the maximum number of bytes available for digest value storage. However, the display length of statement digests may be longer than the available buffer size due to encoding of statement components such as keywords and literal values in digest buffer. Consequently, values selected from the `DIGEST_TEXT` column of statement event tables may appear to exceed the `performance_schema_max_digest_length` value.

- `CURRENT_SCHEMA`

  The default database for the statement, `NULL` if there is none.

- `OBJECT_SCHEMA`, `OBJECT_NAME`, `OBJECT_TYPE`

  For nested statements (stored programs), these columns contain information about the parent statement. Otherwise they are `NULL`.

- `OBJECT_INSTANCE_BEGIN`

  This column identifies the statement. The value is the address of an object in memory.

- `MYSQL_ERRNO`

  The statement error number, from the statement diagnostics area.

- `RETURNED_SQLSTATE`

  The statement SQLSTATE value, from the statement diagnostics area.

- `MESSAGE_TEXT`

  The statement error message, from the statement diagnostics area.

- `ERRORS`

  Whether an error occurred for the statement. The value is 0 if the SQLSTATE value begins with `00` (completion) or `01` (warning). The value is 1 is the SQLSTATE value is anything else.

- `WARNINGS`

  The number of warnings, from the statement diagnostics area.

- `ROWS_AFFECTED`

  The number of rows affected by the statement. For a description of the meaning of “affected,” see Section 27.7.7.1, “mysql_affected_rows()”.

- `ROWS_SENT`

  The number of rows returned by the statement.

- `ROWS_EXAMINED`

  The number of rows read from storage engines during statement execution.

- `CREATED_TMP_DISK_TABLES`

  Like the `Created_tmp_disk_tables` status variable, but specific to the statement.

- `CREATED_TMP_TABLES`

  Like the `Created_tmp_tables` status variable, but specific to the statement.

- `SELECT_FULL_JOIN`

  Like the `Select_full_join` status variable, but specific to the statement.

- `SELECT_FULL_RANGE_JOIN`

  Like the `Select_full_range_join` status variable, but specific to the statement.

- `SELECT_RANGE`

  Like the `Select_range` status variable, but specific to the statement.

- `SELECT_RANGE_CHECK`

  Like the `Select_range_check` status variable, but specific to the statement.

- `SELECT_SCAN`

  Like the `Select_scan` status variable, but specific to the statement.

- `SORT_MERGE_PASSES`

  Like the `Sort_merge_passes` status variable, but specific to the statement.

- `SORT_RANGE`

  Like the `Sort_range` status variable, but specific to the statement.

- `SORT_ROWS`

  Like the `Sort_rows` status variable, but specific to the statement.

- `SORT_SCAN`

  Like the `Sort_scan` status variable, but specific to the statement.

- `NO_INDEX_USED`

  1 if the statement performed a table scan without using an index, 0 otherwise.

- `NO_GOOD_INDEX_USED`

  1 if the server found no good index to use for the statement, 0 otherwise. For additional information, see the description of the `Extra` column from `EXPLAIN` output for the `Range checked for each record` value in Section 8.8.2, “EXPLAIN Output Format”.

- `NESTING_EVENT_ID`, `NESTING_EVENT_TYPE`, `NESTING_EVENT_LEVEL`

  These three columns are used with other columns to provide information as follows for top-level (unnested) statements and nested statements (executed within a stored program).

  For top level statements:

```
OBJECT_TYPE = NULL
OBJECT_SCHEMA = NULL
OBJECT_NAME = NULL
NESTING_EVENT_ID = NULL
NESTING_EVENT_TYPE = NULL
NESTING_LEVEL = 0
```

  For nested statements:

```
OBJECT_TYPE = the parent statement object type
OBJECT_SCHEMA = the parent statement object schema
OBJECT_NAME = the parent statement object name
NESTING_EVENT_ID = the parent statement EVENT_ID
NESTING_EVENT_TYPE = 'STATEMENT'
NESTING_LEVEL = the parent statement NESTING_LEVEL plus one
```

The `events_statements_current` table has these indexes:

- Primary key on (`THREAD_ID`, `EVENT_ID`)

`TRUNCATE TABLE` is permitted for the `events_statements_current` table. It removes the rows.

25.11.6.2 The events_statements_history Table

The `events_statements_history` table contains the most recent *`N`* statement events per thread. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_statements_history_size` system variable at server startup. Statement events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full.

The `events_statements_history` table has the same structure as `events_statements_current`, including indexing. See Section 25.11.6.1, “The events_statements_current Table”.

`TRUNCATE TABLE` is permitted for the `events_statements_history` table. It removes the rows.

For information about configuration of statement event collection, see Section 25.11.6, “Performance Schema Statement Event Tables”.

25.11.6.3 The events_statements_history_long Table

The `events_statements_history_long` table contains the most recent *`N`* statement events. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_statements_history_long_size` system variable at server startup. Statement events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full. When a thread ends, its rows are removed from the table.

The `events_statements_history_long` table has the same structure as `events_statements_current`, except that it has no indexes. See Section 25.11.6.1, "The events_statements_current Table".

`TRUNCATE TABLE` is permitted for the `events_statements_history_long` table. It removes the rows.

For information about configuration of statement event collection, see Section 25.11.6, "Performance Schema Statement Event Tables".

25.11.6.4 The prepared_statements_instances Table

The Performance Schema provides instrumentation for prepared statements, for which there are two protocols:

- The binary protocol. This is accessed through the MySQL C API and maps onto underlying server commands as shown in the following table.

C API Function	Corresponding Server Command
`mysql_stmt_prepare()`	`COM_STMT_PREPARE`
`mysql_stmt_execute()`	`COM_STMT_EXECUTE`
`mysql_stmt_close()`	`COM_STMT_CLOSE`

- The text protocol. This is accessed using SQL statements and maps onto underlying server commands as shown in the following table.

SQL Statement	Corresponding Server Command
`PREPARE`	`SQLCOM_PREPARE`
`EXECUTE`	`SQLCOM_EXECUTE`
`DEALLOCATE PREPARE`, `DROP PREPARE`	`SQLCOM_DEALLOCATE PREPARE`

Performance Schema prepared statement instrumentation covers both protocols. The following discussion refers to the server commands rather than the C API functions or SQL statements.

Information about prepared statements is available in the `prepared_statements_instances` table. This table enables inspection of prepared statements used in the server and provides aggregated statistics about them. To control the size of this table, set the `performance_schema_max_prepared_statements_instances` system variable at server startup.

Collection of prepared statement information depends on the statement instruments shown in the following table. These instruments are enabled by default. To modify them, update the `setup_instruments` table.

Instrument	Server Command
`statement/com/Prepare`	`COM_STMT_PREPARE`
`statement/com/Execute`	`COM_STMT_EXECUTE`
`statement/sql/prepare_sql`	`SQLCOM_PREPARE`
`statement/sql/execute_sql`	`SQLCOM_EXECUTE`

The Performance Schema manages the contents of the `prepared_statements_instances` table as follows:

- Statement preparation

  A `COM_STMT_PREPARE` or `SQLCOM_PREPARE` command creates a prepared statement in the server. If the statement is successfully instrumented, a new row is added to the `prepared_statements_instances` table. If the statement cannot be instrumented, `Performance_schema_prepared_statements_lost` status variable is incremented.

- Prepared statement execution

  Execution of a `COM_STMT_EXECUTE` or `SQLCOM_PREPARE` command for an instrumented prepared statement instance updates the corresponding `prepared_statements_instances` table row.

- Prepared statement deallocation

  Execution of a `COM_STMT_CLOSE` or `SQLCOM_DEALLOCATE_PREPARE` command for an instrumented prepared statement instance removes the corresponding `prepared_statements_instances` table row. To avoid resource leaks, removal occurs even if the prepared statement instruments described previously are disabled.

The `prepared_statements_instances` table has these columns:

- `OBJECT_INSTANCE_BEGIN`

  The address in memory of the instrumented prepared statement.

- `STATEMENT_ID`

  The internal statement ID assigned by the server. The text and binary protocols both use statement IDs.

- `STATEMENT_NAME`

  For the binary protocol, this column is `NULL`. For the text protocol, this column is the external statement name assigned by the user. For example, for the following SQL statement, the name of the prepared statement is `stmt`:

  ```
  PREPARE stmt FROM 'SELECT 1';
  ```

- `SQL_TEXT`

  The prepared statement text, with `?` placeholder markers.

- `OWNER_THREAD_ID`, `OWNER_EVENT_ID`

  These columns indicate the event that created the prepared statement.

- `OWNER_OBJECT_TYPE`, `OWNER_OBJECT_SCHEMA`, `OWNER_OBJECT_NAME`

  For a prepared statement created by a client session, these columns are `NULL`. For a prepared statement created by a stored program, these columns point to the stored program. A typical user error is forgetting to deallocate prepared statements. These columns can be used to find stored programs that leak prepared statements:

  ```
  SELECT OWNER_OBJECT_TYPE, OWNER_OBJECT_SCHEMA, OWNER_OBJECT_NAME,
  STATEMENT_NAME, SQL_TEXT
  FROM performance_schema.prepared_statements_instances
  WHERE OWNER_OBJECT_TYPE IS NOT NULL;
  ```

- `TIMER_PREPARE`

The time spent executing the statement preparation itself.

- `COUNT_REPREPARE`

  The number of times the statement was reprepared internally (see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”). Timing statistics for repreparation are not available because it is counted as part of statement execution, not as a separate operation.

- `COUNT_EXECUTE`, `SUM_TIMER_EXECUTE`, `MIN_TIMER_EXECUTE`, `AVG_TIMER_EXECUTE`, `MAX_TIMER_EXECUTE`

  Aggregated statistics for executions of the prepared statement.

- `SUM_xxx`

  The remaining `SUM_xxx` columns are the same as for the statement summary tables (see Section 25.11.15.3, “Statement Summary Tables”).

The `prepared_statements_instances` table has these indexes:

- Primary key on (`OBJECT_INSTANCE_BEGIN`)
- Index on (`STATEMENT_ID`)
- Index on (`STATEMENT_NAME`)
- Index on (`OWNER_THREAD_ID`, `OWNER_EVENT_ID`)
- Index on (`OWNER_OBJECT_TYPE`, `OWNER_OBJECT_SCHEMA`, `OWNER_OBJECT_NAME`)

`TRUNCATE TABLE` resets the statistics columns of the `prepared_statements_instances` table.

25.11.7 Performance Schema Transaction Tables

The Performance Schema instruments transactions. Within the event hierarchy, wait events nest within stage events, which nest within statement events, which nest within transaction events.

These tables store transaction events:

- `events_transactions_current`: Current transaction events
- `events_transactions_history`: The most recent transaction events per thread
- `events_transactions_history_long`: The most recent transaction events globally (across all threads)

The following sections describe the transaction event tables. There are also summary tables that aggregate information about transaction events; see Section 25.11.15.5, “Transaction Summary Tables”.

Transaction Event Configuration

To control collection of transaction events, set the state of the relevant instruments and consumers:

- The `setup_instruments` table contains an instrument named `transaction`. Use this instrument to enable or disable collection of transaction events.
- The `setup_consumers` table contains consumer values with names corresponding to the current and recent transaction event table names. Use these consumers to filter collection of transaction events.

The `transaction` instrument and the `events_transactions_current` and `events_transactions_history` transaction consumers are enabled by default:

```
mysql> SELECT * FROM setup_instruments WHERE NAME = 'transaction';
+-------------+---------+-------+
| NAME        | ENABLED | TIMED |
+-------------+---------+-------+
| transaction | YES     | YES   |
+-------------+---------+-------+
mysql> SELECT * FROM setup_consumers WHERE NAME LIKE '%transactions%';
+---------------------------------+---------+
| NAME                            | ENABLED |
+---------------------------------+---------+
| events_transactions_current     | YES     |
| events_transactions_history     | YES     |
| events_transactions_history_long | NO      |
+---------------------------------+---------+
```

(Prior to MySQL 8.0.2, the `transaction` instrument and the transaction consumers are disabled by default.)

To control transaction event collection at server startup, use lines like these in your `my.cnf` file:

- Enable:

```
[mysqld]
performance-schema-instrument='transaction=ON'
performance-schema-consumer-events-transactions-current=ON
performance-schema-consumer-events-transactions-history=ON
performance-schema-consumer-events-transactions-history-long=ON
```

- Disable:

```
[mysqld]
performance-schema-instrument='transaction=OFF'
performance-schema-consumer-events-transactions-current=OFF
performance-schema-consumer-events-transactions-history=OFF
performance-schema-consumer-events-transactions-history-long=OFF
```

To control transaction event collection at runtime, update the `setup_instruments` and `setup_consumers` tables:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
WHERE NAME = 'transaction';
UPDATE setup_consumers SET ENABLED = 'YES'
WHERE NAME LIKE '%transactions%';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO', TIMED = 'NO'
WHERE NAME = 'transaction';
UPDATE setup_consumers SET ENABLED = 'NO'
WHERE NAME LIKE '%transactions%';
```

To collect transaction events only for specific transaction event tables, enable the `transaction` instrument but only the transaction consumers corresponding to the desired tables.

The `setup_timers` table contains a row with a `NAME` value of `transaction` that indicates the unit for transaction event timing. The default unit is `NANOSECOND`:

```
mysql> SELECT * FROM setup_timers WHERE NAME = 'transaction';
+-------------+------------+
| NAME        | TIMER_NAME |
+-------------+------------+
| transaction | NANOSECOND |
+-------------+------------+
```

To change the timing unit, modify the `TIMER_NAME` value:

```
UPDATE setup_timers SET TIMER_NAME = 'MICROSECOND'
WHERE NAME = 'transaction';
```

For additional information about configuring event collection, see Section 25.3, “Performance Schema Startup Configuration”, and Section 25.4, “Performance Schema Runtime Configuration”.

Transaction Boundaries

In MySQL Server, transactions start explicitly with these statements:

```
START TRANSACTION | BEGIN | XA START | XA BEGIN
```

Transactions also start implicitly. For example, when the `autocommit` system variable is enabled, the start of each statement starts a new transaction.

When `autocommit` is disabled, the first statement following a committed transaction marks the start of a new transaction. Subsequent statements are part of the transaction until it is committed.

Transactions explicitly end with these statements:

```
COMMIT | ROLLBACK | XA COMMIT | XA ROLLBACK
```

Transactions also end implicitly, by execution of DDL statements, locking statements, and server administration statements.

In the following discussion, references to `START TRANSACTION` also apply to `BEGIN`, `XA START`, and `XA BEGIN`. Similarly, references to `COMMIT` and `ROLLBACK` apply to `XA COMMIT` and `XA ROLLBACK`, respectively.

The Performance Schema defines transaction boundaries similarly to that of the server. The start and end of a transaction event closely match the corresponding state transitions in the server:

- For an explicitly started transaction, the transaction event starts during processing of the `START TRANSACTION` statement.
- For an implicitly started transaction, the transaction event starts on the first statement that uses a transactional engine after the previous transaction has ended.
- For any transaction, whether explicitly or implicitly ended, the transaction event ends when the server transitions out of the active transaction state during the processing of `COMMIT` or `ROLLBACK`.

There are subtle implications to this approach:

- Transaction events in the Performance Schema do not fully include the statement events associated with the corresponding `START TRANSACTION`, `COMMIT`, or `ROLLBACK` statements. There is a trivial amount of timing overlap between the transaction event and these statements.
- Statements that work with nontransactional engines have no effect on the transaction state of the connection. For implicit transactions, the transaction event begins with the first statement that uses a transactional engine. This means that statements operating exclusively on nontransactional tables are ignored, even following `START TRANSACTION`.

To illustrate, consider the following scenario:

```
1. SET autocommit = OFF;
2. CREATE TABLE t1 (a INT) ENGINE = InnoDB;
3. START TRANSACTION;                       -- Transaction 1 START
4. INSERT INTO t1 VALUES (1), (2), (3);
5. CREATE TABLE t2 (a INT) ENGINE = MyISAM; -- Transaction 1 COMMIT
                                            -- (implicit; DDL forces commit)
6. INSERT INTO t2 VALUES (1), (2), (3);     -- Update nontransactional table
7. UPDATE t2 SET a = a + 1;                 -- ... and again
8. INSERT INTO t1 VALUES (4), (5), (6);     -- Write to transactional table
                                            -- Transaction 2 START (implicit)
9. COMMIT;                                  -- Transaction 2 COMMIT
```

From the perspective of the server, Transaction 1 ends when table `t2` is created. Transaction 2 does not start until a transactional table is accessed, despite the intervening updates to nontransactional tables.

From the perspective of the Performance Schema, Transaction 2 starts when the server transitions into an active transaction state. Statements 6 and 7 are not included within the boundaries of Transaction 2, which is consistent with how the server writes transactions to the binary log.

Transaction Instrumentation

Three attributes define transactions:

- Access mode (read only, read write)
- Isolation level (`SERIALIZABLE`, `REPEATABLE READ`, and so forth)
- Implicit (`autocommit` enabled) or explicit (`autocommit` disabled)

To reduce complexity of the transaction instrumentation and to ensure that the collected transaction data provides complete, meaningful results, all transactions are instrumented independently of access mode, isolation level, or autocommit mode.

To selectively examine transaction history, use the attribute columns in the transaction event tables: `ACCESS_MODE`, `ISOLATION_LEVEL`, and `AUTOCOMMIT`.

The cost of transaction instrumentation can be reduced various ways, such as enabling or disabling transaction instrumentation according to user, account, host, or thread (client connection).

Transactions and Nested Events

The parent of a transaction event is the event that initiated the transaction. For an explicitly started transaction, this includes the `START TRANSACTION` and `COMMIT AND CHAIN` statements. For an implicitly started transaction, it is the first statement that uses a transactional engine after the previous transaction ends.

In general, a transaction is the top-level parent to all events initiated during the transaction, including statements that explicitly end the transaction such as `COMMIT` and `ROLLBACK`. Exceptions are statements that implicitly end a transaction, such as DDL statements, in which case the current transaction must be committed before the new statement is executed.

Transactions and Stored Programs

Transactions and stored program events are related as follows:

- Stored Procedures

  Stored procedures operate independently of transactions. A stored procedure can be started within a transaction, and a transaction can be started or ended from within a stored procedure. If called from

within a transaction, a stored procedure can execute statements that force a commit of the parent transaction and then start a new transaction.

If a stored procedure is started within a transaction, that transaction is the parent of the stored procedure event.

If a transaction is started by a stored procedure, the stored procedure is the parent of the transaction event.

- Stored Functions

  Stored functions are restricted from causing an explicit or implicit commit or rollback. Stored function events can reside within a parent transaction event.

- Triggers

  Triggers activate as part of a statement that accesses the table with which it is associated, so the parent of a trigger event is always the statement that activates it.

  Triggers cannot issue statements that cause an explicit or implicit commit or rollback of a transaction.

- Scheduled Events

  The execution of the statements in the body of a scheduled event takes place in a new connection. Nesting of a scheduled event within a parent transaction is not applicable.

Transactions and Savepoints

Savepoint statements are recorded as separate statement events. Transaction events include separate counters for `SAVEPOINT`, `ROLLBACK TO SAVEPOINT`, and `RELEASE SAVEPOINT` statements issued during the transaction.

Transactions and Errors

Errors and warnings that occur within a transaction are recorded in statement events, but not in the corresponding transaction event. This includes transaction-specific errors and warnings, such as a rollback on a nontransactional table or GTID consistency errors.

25.11.7.1 The events_transactions_current Table

The `events_transactions_current` table contains current transaction events, one row per thread showing the current status of the thread's most recent monitored transaction event. For example:

```
mysql> SELECT * FROM events_transactions_current LIMIT 1\G
*************************** 1. row ***************************
                    THREAD_ID: 26
                     EVENT_ID: 7
                 END_EVENT_ID: NULL
                   EVENT_NAME: transaction
                        STATE: ACTIVE
                       TRX_ID: NULL
                         GTID: 3E11FA47-71CA-11E1-9E33-C80AA9429562:56
                          XID: NULL
                     XA_STATE: NULL
                       SOURCE: transaction.cc:150
                  TIMER_START: 420833537900000
                    TIMER_END: NULL
                   TIMER_WAIT: NULL
                  ACCESS_MODE: READ WRITE
              ISOLATION_LEVEL: REPEATABLE READ
                   AUTOCOMMIT: NO
           NUMBER_OF_SAVEPOINTS: 0
NUMBER_OF_ROLLBACK_TO_SAVEPOINT: 0
    NUMBER_OF_RELEASE_SAVEPOINT: 0
```

```
          OBJECT_INSTANCE_BEGIN: NULL
               NESTING_EVENT_ID: 6
             NESTING_EVENT_TYPE: STATEMENT
```

Of the tables that contain transaction event rows, `events_transactions_current` is the most fundamental. Other tables that contain transaction event rows are logically derived from the current events. For example, the `events_transactions_history` and `events_transactions_history_long` tables are collections of the most recent transaction events, up to a fixed number of rows.

For information about configuration of transaction event collection, see Section 25.11.7, “Performance Schema Transaction Tables”.

The `events_transactions_current` table has these columns:

- `THREAD_ID`, `EVENT_ID`

  The thread associated with the event and the thread current event number when the event starts. The `THREAD_ID` and `EVENT_ID` values taken together uniquely identify the row. No two rows have the same pair of values.

- `END_EVENT_ID`

  This column is set to `NULL` when the event starts and updated to the thread current event number when the event ends.

- `EVENT_NAME`

  The name of the instrument from which the event was collected. This is a `NAME` value from the `setup_instruments` table. Instrument names may have multiple parts and form a hierarchy, as discussed in Section 25.6, “Performance Schema Instrument Naming Conventions”.

- `STATE`

  The current transaction state. The value is `ACTIVE` (after `START TRANSACTION` or `BEGIN`), `COMMITTED` (after `COMMIT`), or `ROLLED BACK` (after `ROLLBACK`).

- `TRX_ID`

  Unused.

- `GTID`

  The GTID column contains the value of `gtid_next`, which can be one of `ANONYMOUS`, `AUTOMATIC`, or a GTID using the format `UUID:NUMBER`. For transactions that use `gtid_next=AUTOMATIC`, which is all normal client transactions, the GTID column changes when the transaction commits and the actual GTID is assigned. If `gtid_mode` is either `ON` or `ON_PERMISSIVE`, the GTID column changes to the transaction's GTID. If `gtid_mode` is either `OFF` or `OFF_PERMISSIVE`, the GTID column changes to `ANONYMOUS`.

- `XID_FORMAT_ID`, `XID_GTRID`, and `XID_BQUAL`

  The components of the XA transaction identifier. They have the format described in Section 13.3.7.1, “XA Transaction SQL Syntax”.

- `XA_STATE`

  The state of the XA transaction. The value is `ACTIVE` (after `XA START`), `IDLE` (after `XA END`), `PREPARED` (after `XA PREPARE`), `ROLLED BACK` (after `XA ROLLBACK`), or `COMMITTED` (after `XA COMMIT`).

- `SOURCE`

The name of the source file containing the instrumented code that produced the event and the line number in the file at which the instrumentation occurs. This enables you to check the source to determine exactly what code is involved.

- `TIMER_START`, `TIMER_END`, `TIMER_WAIT`

  Timing information for the event. The unit for these values is picoseconds (trillionths of a second). The `TIMER_START` and `TIMER_END` values indicate when event timing started and ended. `TIMER_WAIT` is the event elapsed time (duration).

  If an event has not finished, `TIMER_END` is the current timer value and `TIMER_WAIT` is the time elapsed so far (`TIMER_END` − `TIMER_START`).

  If an event is produced from an instrument that has `TIMED = NO`, timing information is not collected, and `TIMER_START`, `TIMER_END`, and `TIMER_WAIT` are all `NULL`.

  For discussion of picoseconds as the unit for event times and factors that affect time values, see Section 25.4.1, “Performance Schema Event Timing”.

- `ACCESS_MODE`

  The transaction access mode. The value is `READ ONLY` or `READ WRITE`.

- `ISOLATION_LEVEL`

  The transaction isolation level. The value is `REPEATABLE READ`, `READ COMMITTED`, `READ UNCOMMITTED`, or `SERIALIZABLE`.

- `AUTOCOMMIT`

  Whether autcommit mode was enabled when the transaction started.

- `NUMBER_OF_SAVEPOINTS`, `NUMBER_OF_ROLLBACK_TO_SAVEPOINT`, `NUMBER_OF_RELEASE_SAVEPOINT`

  The number of `SAVEPOINT`, `ROLLBACK TO SAVEPOINT`, and `RELEASE SAVEPOINT` statements issued during the transaction.

- `OBJECT_INSTANCE_BEGIN`

  Unused.

- `NESTING_EVENT_ID`

  The `EVENT_ID` value of the event within which this event is nested.

- `NESTING_EVENT_TYPE`

  The nesting event type. The value is `TRANSACTION`, `STATEMENT`, `STAGE`, or `WAIT`. (`TRANSACTION` will not appear because transactions cannot be nested.)

The `events_transactions_current` table has these indexes:

- Primary key on (`THREAD_ID`, `EVENT_ID`)

`TRUNCATE TABLE` is permitted for the `events_transactions_current` table. It removes the rows.

25.11.7.2 The events_transactions_history Table

The `events_transactions_history` table contains the most recent *`N`* transaction events per thread. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_transactions_history_size` system variable at server startup.

Transaction events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full.

The `events_transactions_history` table has the same structure as `events_transactions_current`, including indexing. See Section 25.11.7.1, "The events_transactions_current Table".

`TRUNCATE TABLE` is permitted for the `events_transactions_history` table. It removes the rows.

For information about configuration of transaction event collection, see Section 25.11.7, "Performance Schema Transaction Tables".

25.11.7.3 The events_transactions_history_long Table

The `events_transactions_history_long` table contains the most recent *`N`* transaction events. The value of *`N`* is autosized at server startup. To set the table size explicitly, set the `performance_schema_events_transactions_history_long_size` system variable at server startup. Transaction events are not added to the table until they have ended. As new events are added, older events are discarded if the table is full. When a thread ends, its rows are removed from the table.

The `events_transactions_history_long` table has the same structure as `events_transactions_current`, except that it has no indexes. See Section 25.11.7.1, "The events_transactions_current Table".

`TRUNCATE TABLE` is permitted for the `events_transactions_history_long` table. It removes the rows.

For information about configuration of transaction event collection, see Section 25.11.7, "Performance Schema Transaction Tables".

25.11.8 Performance Schema Connection Tables

When a client connects to the MySQL server, it does so under a particular user name and from a particular host. The Performance Schema provides statistics about these connections, tracking them per account (user and host combination) as well as separately per user name and host name, using these tables:

- `accounts`: Connection statistics per client account
- `hosts`: Connection statistics per client host name
- `users`: Connection statistics per client user name

The meaning of "account" in the connection tables is similar to its meaning in the MySQL grant tables in the `mysql` system database, in the sense that the term refers to a combination of user and host values. They differ in that, for grant tables, the host part of an account can be a pattern, whereas for Performance Schema tables, the host value is always a specific nonpattern host name.

Each connection table has `CURRENT_CONNECTIONS` and `TOTAL_CONNECTIONS` columns to track the current and total number of connections per "tracking value" on which its statistics are based. The tables differ in what they use for the tracking value. The `accounts` table has `USER` and `HOST` columns to track connections per user and host combination. The `users` and `hosts` tables have a `USER` and `HOST` column, respectively, to track connections per user name and host name.

The Performance Schema also counts internal threads and threads for user sessions that failed to authenticate, using rows with `USER` and `HOST` column values of `NULL`.

Suppose that clients named `user1` and `user2` each connect one time from `hosta` and `hostb`. The Performance Schema tracks the connections as follows:

- The `accounts` table has four rows, for the `user1`/`hosta`, `user1`/`hostb`, `user2`/`hosta`, and `user2`/`hostb` account values, each row counting one connection per account.

- The `hosts` table has two rows, for `hosta` and `hostb`, each row counting two connections per host name.
- The `users` table has two rows, for `user1` and `user2`, each row counting two connections per user name.

When a client connects, the Performance Schema determines which row in each connection table applies, using the tracking value appropriate to each table. If there is no such row, one is added. Then the Performance Schema increments by one the `CURRENT_CONNECTIONS` and `TOTAL_CONNECTIONS` columns in that row.

When a client disconnects, the Performance Schema decrements by one the `CURRENT_CONNECTIONS` column in the row and leaves the `TOTAL_CONNECTIONS` column unchanged.

`TRUNCATE TABLE` is permitted for connection tables. It has these effects:

- Rows are removed for accounts, hosts, or users that have no current connections (rows with `CURRENT_CONNECTIONS = 0`).
- Nonremoved rows are reset to count only current connections: For rows with `CURRENT_CONNECTIONS > 0`, `TOTAL_CONNECTIONS` is reset to `CURRENT_CONNECTIONS`.
- Summary tables that depend on the connection table are implicitly truncated, as described later in this section.

The Performance Schema maintains summary tables that aggregate connection statistics for various event types by account, host, or user. These tables have `_summary_by_account`, `_summary_by_host`, or `_summary_by_user` in the name. To identify them, use this query:

```
mysql> SELECT TABLE_NAME FROM INFORMATION_SCHEMA.TABLES
       WHERE TABLE_SCHEMA = 'performance_schema'
       AND TABLE_NAME REGEXP '_summary_by_(account|host|user)'
       ORDER BY TABLE_NAME;
+-----------------------------------------------------+
| TABLE_NAME                                          |
+-----------------------------------------------------+
| events_errors_summary_by_account_by_error           |
| events_errors_summary_by_host_by_error              |
| events_errors_summary_by_user_by_error              |
| events_stages_summary_by_account_by_event_name      |
| events_stages_summary_by_host_by_event_name         |
| events_stages_summary_by_user_by_event_name         |
| events_statements_summary_by_account_by_event_name  |
| events_statements_summary_by_host_by_event_name     |
| events_statements_summary_by_user_by_event_name     |
| events_transactions_summary_by_account_by_event_name |
| events_transactions_summary_by_host_by_event_name   |
| events_transactions_summary_by_user_by_event_name   |
| events_waits_summary_by_account_by_event_name       |
| events_waits_summary_by_host_by_event_name          |
| events_waits_summary_by_user_by_event_name          |
| memory_summary_by_account_by_event_name             |
| memory_summary_by_host_by_event_name                |
| memory_summary_by_user_by_event_name                |
+-----------------------------------------------------+
```

For details about individual connection summary tables, consult the section that describes tables for the summarized event type:

- Wait event summaries: Section 25.11.15.1, “Wait Event Summary Tables”
- Stage event summaries: Section 25.11.15.2, “Stage Summary Tables”
- Statement event summaries: Section 25.11.15.3, “Statement Summary Tables”
- Transaction event summaries: Section 25.11.15.5, “Transaction Summary Tables”

- Memory event summaries: Section 25.11.15.10, “Memory Summary Tables”
- Error event summaries: Section 25.11.15.11, “Error Summary Tables”

`TRUNCATE TABLE` is permitted for connection summary tables. It removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero for the remaining rows. In addition, each summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends. The following table describes the relationship between connection table truncation and implicitly truncated tables.

Table 25.2 Implicit Effects of Connection Table Truncation

Truncated Connection Table	Implicitly Truncated Summary Tables
`accounts`	Tables with names containing `_summary_by_account`, `_summary_by_thread`
`hosts`	Tables with names containing `_summary_by_account`, `_summary_by_host`, `_summary_by_thread`
`users`	Tables with names containing `_summary_by_account`, `_summary_by_user`, `_summary_by_thread`

Truncating a `_summary_global` summary table also implicitly truncates its corresponding connection and thread summary tables. For example, truncating `events_waits_summary_global_by_event_name` implicitly truncates the wait event summary tables that are aggregated by account, host, user, or thread.

25.11.8.1 The accounts Table

The `accounts` table contains a row for each account that has connected to the MySQL server. For each account, the table counts the current and total number of connections. The table size is autosized at server startup. To set the table size explicitly, set the `performance_schema_accounts_size` system variable at server startup. To disable account statistics, set this variable to 0.

The `accounts` table has the following columns. For a description of how the Performance Schema maintains rows in this table, including the effect of `TRUNCATE TABLE`, see Section 25.11.8, “Performance Schema Connection Tables”.

- `USER`

  The client user name for the connection. This is `NULL` for an internal thread, or for a user session that failed to authenticate.

- `HOST`

  The host from which the client connected. This is `NULL` for an internal thread, or for a user session that failed to authenticate.

- `CURRENT_CONNECTIONS`

  The current number of connections for the account.

- `TOTAL_CONNECTIONS`

  The total number of connections for the account.

The `accounts` table has these indexes:

- Primary key on (`USER`, `HOST`)

25.11.8.2 The hosts Table

The `hosts` table contains a row for each host from which clients have connected to the MySQL server. For each host name, the table counts the current and total number of connections. The table size is autosized at server startup. To set the table size explicitly, set the `performance_schema_hosts_size` system variable at server startup. To disable host statistics, set this variable to 0.

The `hosts` table has the following columns. For a description of how the Performance Schema maintains rows in this table, including the effect of `TRUNCATE TABLE`, see Section 25.11.8, "Performance Schema Connection Tables".

- `HOST`

  The host from which the client connected. This is `NULL` for an internal thread, or for a user session that failed to authenticate.

- `CURRENT_CONNECTIONS`

  The current number of connections for the host.

- `TOTAL_CONNECTIONS`

  The total number of connections for the host.

The `hosts` table has these indexes:

- Primary key on (`HOST`)

25.11.8.3 The users Table

The `users` table contains a row for each user who has connected to the MySQL server. For each user name, the table counts the current and total number of connections. The table size is autosized at server startup. To set the table size explicitly, set the `performance_schema_users_size` system variable at server startup. To disable user statistics, set this variable to 0.

The `users` table has the following columns. For a description of how the Performance Schema maintains rows in this table, including the effect of `TRUNCATE TABLE`, see Section 25.11.8, "Performance Schema Connection Tables".

- `USER`

  The client user name for the connection. This is `NULL` for an internal thread, or for a user session that failed to authenticate.

- `CURRENT_CONNECTIONS`

  The current number of connections for the user.

- `TOTAL_CONNECTIONS`

  The total number of connections for the user.

The `users` table has these indexes:

- Primary key on (`USER`)

25.11.9 Performance Schema Connection Attribute Tables

Application programs can provide key/value pairs as connection attributes to be passed to the server at connect time. For the C API, define the attribute set using the `mysql_options()` and `mysql_options4()` functions. Other MySQL Connectors may provide their own attribute-definition methods.

These tables expose attribute information:

- `session_account_connect_attrs`: Connection attributes for the current session, and other sessions associated with the session account
- `session_connect_attrs`: Connection attributes for all sessions

Attribute names that begin with an underscore (_) are reserved for internal use and should not be created by application programs. This convention permits new attributes to be introduced by MySQL without colliding with application attributes.

The set of connection attributes visible on a given connection varies depending on your platform and MySQL Connector used to establish the connection.

The `libmysqlclient` client library (provided in MySQL and MySQL Connector/C distributions) sets these attributes:

- `_client_name`: The client name (`libmysql` for the client library)
- `_client_version`: The client library version
- `_os`: The operating system (for example, `Linux`, `Win64`)
- `_pid`: The client process ID
- `_platform`: The machine platform (for example, `x86_64`)
- `_thread`: The client thread ID (Windows only)

Other MySQL Connectors may define their own connection attributes.

MySQL Connector/J defines these attributes:

- `_client_license`: The connector license type
- `_runtime_vendor`: The Java runtime environment (JRE) vendor
- `_runtime_version`: The Java runtime environment (JRE) version

MySQL Connector/Net defines these attributes:

- `_client_version`: The client library version
- `_os`: The operating system (for example, `Linux`, `Win64`)
- `_pid`: The client process ID
- `_platform`: The machine platform (for example, `x86_64`)
- `_program_name`: The client name
- `_thread`: The client thread ID (Windows only)

PHP defines attributes that depend on how it was compiled:

- Compiled using `libmysqlclient`: The standard `libmysqlclient` attributes, described previously
- Compiled using `mysqlnd`: Only the `_client_name` attribute, with a value of `mysqlnd`

Many MySQL client programs set a `program_name` attribute with a value equal to the client name. For example, `mysqladmin` and `mysqldump` set `program_name` to `mysqladmin` and `mysqldump`, respectively.

Some MySQL client programs define additional attributes:

- `mysqlbinlog` defines the `_client_role` attribute as `binary_log_listener`.
- Replication slave connections define `program_name` as `mysqld`, `_client_role` as `binary_log_listener`, and `_client_replication_channel_name` as the channel name.
- `FEDERATED` storage engine connections define `program_name` as `mysqld` and `_client_role` as `federated_storage`.

There are limits on the amount of connection attribute data transmitted from client to server: A fixed limit imposed by the client prior to connect time; a fixed limit imposed by the server at connect time; and a configurable limit imposed by the Performance Schema at connect time.

For connections initiated using the C API, the `libmysqlclient` library imposes a limit of 64KB on the aggregate size of connection attribute data on the client side: Calls to `mysql_options()` that cause this limit to be exceeded produce a `CR_INVALID_PARAMETER_NO` error. Other MySQL Connectors may impose their own client-side limits on how much connection attribute data can be transmitted to the server.

On the server side, these size checks on connection attribute data occur:

- The server imposes a limit of 64KB on the aggregate size of connection attribute data it will accept. If a client attempts to send more than 64KB of attribute data, the server rejects the connection. Otherwise, the server considers the attribute buffer valid and tracks the size of the longest such buffer in the `Performance_schema_session_connect_attrs_longest_seen` status variable.
- For accepted connections, the Performance Schema checks aggregate attribute size against the value of the `performance_schema_session_connect_attrs_size` system variable. If attribute size exceeds this value, these actions take place:
  - The Performance Schema truncates the attribute data and increments the `Performance_schema_session_connect_attrs_lost` status variable, which indicates the number of connections for which attribute truncation occurred.
  - The Performance Schema writes a message to the error log if the `log_warnings` system variable is greater than zero:

    ```
    Connection attributes of length N were truncated (N bytes lost)
    for connection N, user user_name@host_name (as user_name), auth: {yes|no}
    ```

    The information in the warning message is intended to help DBAs identify clients for which attribute truncation occurred.
  - A `_truncated` attribute is added to the session attributes with a value indicating how many bytes were lost, if the attribute buffer has sufficient space. This enables the Performance Schema to expose per-connection truncation information in the connection attribute tables. This information can be examined without having to check the error log.

25.11.9.1 The session_account_connect_attrs Table

Application programs can provide key/value connection attributes to be passed to the server at connect time, using the `mysql_options()` and `mysql_options4()` C API functions. For descriptions of common attributes, see Section 25.11.9, “Performance Schema Connection Attribute Tables”.

The `session_account_connect_attrs` table contains connection attributes only for sessions for your own account. To see connection attributes for all sessions, look in the `session_connect_attrs` table.

The `session_account_connect_attrs` table contains these columns:

- PROCESSLIST_ID

  The connection identifier for the session.

- ATTR_NAME

  The attribute name.

- ATTR_VALUE

  The attribute value.

- ORDINAL_POSITION

  The order in which the attribute was added to the set of connection attributes.

The session_account_connect_attrs table has these indexes:

- Primary key on (PROCESSLIST_ID, ATTR_NAME)

TRUNCATE TABLE is not permitted for the session_account_connect_attrs table.

25.11.9.2 The session_connect_attrs Table

Application programs can provide key/value connection attributes to be passed to the server at connect time, using the mysql_options() and mysql_options4() C API functions. For descriptions of common attributes, see Section 25.11.9, “Performance Schema Connection Attribute Tables”.

The session_connect_attrs table contains connection attributes for all sessions. To see connection attributes only for sessions for your own account, look in the session_account_connect_attrs table.

The session_connect_attrs table contains these columns:

- PROCESSLIST_ID

  The connection identifier for the session.

- ATTR_NAME

  The attribute name.

- ATTR_VALUE

  The attribute value.

- ORDINAL_POSITION

  The order in which the attribute was added to the set of connection attributes.

The session_connect_attrs table has these indexes:

- Primary key on (PROCESSLIST_ID, ATTR_NAME)

TRUNCATE TABLE is not permitted for the session_connect_attrs table.

25.11.10 Performance Schema User Variable Tables

The Performance Schema provides a user_variables_by_thread table that exposes user-defined variables. These are variables defined within a specific session and include a @ character preceding the name; see Section 9.4, “User-Defined Variables”.

The user_variables_by_thread table contains these columns:

- `THREAD_ID`

  The thread identifier of the session in which the variable is defined.

- `VARIABLE_NAME`

  The variable name, without the leading `@` character.

- `VARIABLE_VALUE`

  The variable value.

The `user_variables_by_thread` table has these indexes:

- Primary key on (`THREAD_ID`, `VARIABLE_NAME`)

`TRUNCATE TABLE` is not permitted for the `user_variables_by_thread` table.

25.11.11 Performance Schema Replication Tables

The Performance Schema provides tables that expose replication information. This is similar to the information available from the `SHOW SLAVE STATUS` statement, but representation in table form is more accessible and has usability benefits:

- `SHOW SLAVE STATUS` output is useful for visual inspection, but not so much for programmatic use. By contrast, using the Performance Schema tables, information about slave status can be searched using general `SELECT` queries, including complex `WHERE` conditions, joins, and so forth.
- Query results can be saved in tables for further analysis, or assigned to variables and thus used in stored procedures.
- The replication tables provide better diagnostic information. For multi-threaded slave operation, `SHOW SLAVE STATUS` reports all coordinator and worker thread errors using the `Last_SQL_Errno` and `Last_SQL_Error` fields, so only the most recent of those errors is visible and information can be lost. The replication tables store errors on a per-thread basis without loss of information.
- The last seen transaction is visible in the replication tables on a per-worker basis. This is information not avilable from `SHOW SLAVE STATUS`.
- Developers familiar with the Performance Schema interface can extend the replication tables to provide additional information by adding rows to the tables.

Replication Table Descriptions

The Performance Schema provides the following replication-related tables:

- Tables that contain information about the connection of the slave server to the master server:
  - `replication_connection_configuration`: Configuration parameters for connecting to the master
  - `replication_connection_status`: Current status of the connection to the master
- Tables that contain general (not thread-specific) information about the transaction applier:
  - `replication_applier_configuration`: Configuration parameters for the transaction applier on the slave.
  - `replication_applier_status`: Current status of the transaction applier on the slave.
- Tables that contain information about specific threads responsible for applying transactions received from the master:

  - `replication_applier_status_by_coordinator`: Status of the coordinator thread (empty unless the slave is multi-threaded).
  - `replication_applier_status_by_worker`: Status of the applier thread or worker threads if the slave is multi-threaded.
- Tables that contain information about channel based replication filters:
  - `replication_applier_filters`: Provides information about the replication filters configured on specific replication channels.
  - `replication_applier_global_filters`: Provides information about global replication filters, which apply to all replication channels.
- Tables that contain information about Group Replication members:
  - `replication_group_members`: Provides network and status information for group members.
  - `replication_group_member_stats`: Provides statistical information about group members and transactions in which they participate.

  For more information see Section 21.3, "Monitoring Group Replication".

The following sections describe each replication table in more detail, including the correspondence between the columns produced by `SHOW SLAVE STATUS` and the replication table columns in which the same information appears.

The remainder of this introduction to the replication tables describes how the Performance Schema populates them and which fields from `SHOW SLAVE STATUS` are not represented in the tables.

Replication Table Life Cycle

The Performance Schema populates the replication tables as follows:

- Prior to execution of `CHANGE MASTER TO`, the tables are empty.
- After `CHANGE MASTER TO`, the configuration parameters can be seen in the tables. At this time, there are no active slave threads, so the `THREAD_ID` columns are `NULL` and the `SERVICE_STATE` columns have a value of `OFF`.
- After `START SLAVE`, non-`NULL` `THREAD_ID` values can be seen. Threads that are idle or active have a `SERVICE_STATE` value of `ON`. The thread that connects to the master server has a value of `CONNECTING` while it establishes the connection, and `ON` thereafter as long as the connection lasts.
- After `STOP SLAVE`, the `THREAD_ID` columns become `NULL` and the `SERVICE_STATE` columns for threads that no longer exist have a value of `OFF`.
- The tables are preserved after `STOP SLAVE` or threads dying due to an error.
- The `replication_applier_status_by_worker` table is nonempty only when the slave is operating in multi-threaded mode. That is, if the `slave_parallel_workers` system variable is greater than 0, this table is populated when `START SLAVE` is executed, and the number of rows shows the number of workers.

SHOW SLAVE STATUS Information Not In the Replication Tables

The information in the Performance Schema replication tables differs somewhat from the information available from `SHOW SLAVE STATUS` because the tables are oriented toward use of global transaction identifiers (GTIDs), not file names and positions, and they represent server UUID values, not server ID values. Due to these differences, several `SHOW SLAVE STATUS` columns are not preserved in the Performance Schema replication tables, or are represented a different way:

- The following fields refer to file names and positions and are not preserved:

```
Master_Log_File
Read_Master_Log_Pos
Relay_Log_File
Relay_Log_Pos
Relay_Master_Log_File
Exec_Master_Log_Pos
Until_Condition
Until_Log_File
Until_Log_Pos
```

- The `Master_Info_File` field is not preserved. It refers to the `master.info` file, which has been superseded by crash-safe slave tables.
- The following fields are based on `server_id`, not `server_uuid`, and are not preserved:

```
Master_Server_Id
Replicate_Ignore_Server_Ids
```

- The `Skip_Counter` field is based on event counts, not GTIDs, and is not preserved.
- These error fields are aliases for `Last_SQL_Errno` and `Last_SQL_Error`, so they are not preserved:

```
Last_Errno
Last_Error
```

  In the Performance Schema, this error information is available in the `LAST_ERROR_NUMBER` and `LAST_ERROR_MESSAGE` columns of the `replication_applier_status_by_coordinator` table (and `replication_applier_status_by_worker` if the slave is multi-threaded). Those tables provide more specific per-thread error information than is available from `Last_Errno` and `Last_Error`.

- Fields that provide information about command-line filtering options is not preserved:

```
Replicate_Do_DB
Replicate_Ignore_DB
Replicate_Do_Table
Replicate_Ignore_Table
Replicate_Wild_Do_Table
Replicate_Wild_Ignore_Table
```

- The `Slave_IO_State` and `Slave_SQL_Running_State` fields are not preserved. If needed, these values can be obtained from the process list by using the `THREAD_ID` column of the appropriate replication table and joining it with the `ID` column in the `INFORMATION_SCHEMA PROCESSLIST` table to select the `STATE` column of the latter table.
- The `Executed_Gtid_Set` field can show a large set with a great deal of text. Instead, the Performance Schema tables show GTIDs of transactions that are currently being applied by the slave. Alternatively, the set of executed GTIDs can be obtained from the value of the `gtid_executed` system variable.
- The `Seconds_Behind_Master` and `Relay_Log_Space` fields are in to-be-decided status and are not preserved.

Status Variables Moved to Replication Tables

As of MySQL version 5.7.5, the following status variables (previously monitored using `SHOW STATUS`) were moved to the Perfomance Schema replication tables:

- `Slave_retried_transactions`

- `Slave_last_heartbeat`
- `Slave_received_heartbeats`
- `Slave_heartbeat_period`
- `Slave_running`

These status variables are now only relevant when a single replication channel is being used because they *only* report the status of the default replication channel. When multiple replication channels exist, use the Performance Schema replication tables described in this section, which report these variables for each existing replication channel.

Replication Channels

The first column of the replication Performance Schema tables is `CHANNEL_NAME`. This enables the tables to be viewed per replication channel. In a non-multisource replication setup there is a single default replication channel. When you are using multiple replication channels on a slave, you can filter the tables per replication channel to monitor a specific replication channel. See Section 18.2.3, “Replication Channels” and Section 18.1.4.3, “Multi-Source Replication Monitoring” for more information.

25.11.11.1 The replication_connection_configuration Table

This table shows the configuration parameters used by the slave server for connecting to the master server. Parameters stored in the table can be changed at runtime with the `CHANGE MASTER TO` statement, as indicated in the column descriptions.

Compared to the `replication_connection_status` table, `replication_connection_configuration` changes less frequently. It contains values that define how the slave connects to the master and that remain constant during the connection, whereas `replication_connection_status` contains values that change during the connection.

The `replication_connection_configuration` table has these columns:

- `CHANNEL_NAME`

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

- `HOST`

  The master host that the slave is connected to. (`CHANGE MASTER TO` option: `MASTER_HOST`)

- `PORT`

  The port used to connect to the master. (`CHANGE MASTER TO` option: `MASTER_PORT`)

- `USER`

  The user name of the account used to connect to the master. (`CHANGE MASTER TO` option: `MASTER_USER`)

- `NETWORK_INTERFACE`

  The network interface that the slave is bound to, if any. (`CHANGE MASTER TO` option: `MASTER_BIND`)

- `AUTO_POSITION`

  1 if autopositioning is in use; otherwise 0. (`CHANGE MASTER TO` option: `MASTER_AUTO_POSITION`)

- `SSL_ALLOWED`, `SSL_CA_FILE`, `SSL_CA_PATH`, `SSL_CERTIFICATE`, `SSL_CIPHER`, `SSL_KEY`, `SSL_VERIFY_SERVER_CERTIFICATE`, `SSL_CRL_FILE`, `SSL_CRL_PATH`

  These columns show the SSL parameters used by the slave to connect to the master, if any.

  `SSL_ALLOWED` has these values:

  - `Yes` if an SSL connection to the master is permitted
  - `No` if an SSL connection to the master is not permitted
  - `Ignored` if an SSL connection is permitted but the slave server does not have SSL support enabled

  `CHANGE MASTER TO` options for the other SSL columns: `MASTER_SSL_CA`, `MASTER_SSL_CAPATH`, `MASTER_SSL_CERT`, `MASTER_SSL_CIPHER`, `MASTER_SSL_CRL`, `MASTER_SSL_CRLPATH`, `MASTER_SSL_KEY`, `MASTER_SSL_VERIFY_SERVER_CERT`.

- `CONNECTION_RETRY_INTERVAL`

  The number of seconds between connect retries. (`CHANGE MASTER TO` option: `MASTER_CONNECT_RETRY`)

- `CONNECTION_RETRY_COUNT`

  The number of times the slave can attempt to reconnect to the master in the event of a lost connection. (`CHANGE MASTER TO` option: `MASTER_RETRY_COUNT`)

- `HEARTBEAT_INTERVAL`

  The replication heartbeat interval on a slave, measured in seconds.

The `replication_connection_configuration` table has these indexes:

- Primary key on (`CHANNEL_NAME`)

`TRUNCATE TABLE` is not permitted for the `replication_connection_configuration` table.

The following table shows the correspondence between `replication_connection_configuration` columns and `SHOW SLAVE STATUS` columns.

`replication_connection_configuration` Column	**`SHOW SLAVE STATUS` Column**
`HOST`	`Master_Host`
`PORT`	`Master_Port`
`USER`	`Master_User`
`NETWORK_INTERFACE`	`Master_Bind`
`AUTO_POSITION`	`Auto_Position`
`SSL_ALLOWED`	`Master_SSL_Allowed`
`SSL_CA_FILE`	`Master_SSL_CA_File`
`SSL_CA_PATH`	`Master_SSL_CA_Path`
`SSL_CERTIFICATE`	`Master_SSL_Cert`
`SSL_CIPHER`	`Master_SSL_Cipher`
`SSL_KEY`	`Master_SSL_Key`
`SSL_VERIFY_SERVER_CERTIFICATE`	`Master_SSL_Verify_Server_Cert`
`SSL_CRL_FILE`	`Master_SSL_Crl`
`SSL_CRL_PATH`	`Master_SSL_Crlpath`

replication_connection_configuration Column	SHOW SLAVE STATUS Column
CONNECTION_RETRY_INTERVAL	Connect_Retry
CONNECTION_RETRY_COUNT	Master_Retry_Count

25.11.11.2 The replication_connection_status Table

This table shows the current status of the I/O thread that handles the slave server connection to the master server, information on the last transaction queued in the relay log, and information on the transaction currently being queued in the relay log.

Compared to the `replication_connection_configuration` table, `replication_connection_status` changes more frequently. It contains values that change during the connection, whereas `replication_connection_configuration` contains values which define how the slave connects to the master and that remain constant during the connection.

The `replication_connection_status` table has these columns:

- `CHANNEL_NAME`

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

- `GROUP_NAME`

  If this server is a member of a group, shows the name of the group the server belongs to.

- `SOURCE_UUID`

  The `server_uuid` value from the master.

- `THREAD_ID`

  The I/O thread ID.

- `SERVICE_STATE`

  `ON` (thread exists and is active or idle), `OFF` (thread no longer exists), or `CONNECTING` (thread exists and is connecting to the master).

- `RECEIVED_TRANSACTION_SET`

  The set of global transaction IDs (GTIDs) corresponding to all transactions received by this slave. Empty if GTIDs are not in use. See GTID Sets for more information.

- `LAST_ERROR_NUMBER`, `LAST_ERROR_MESSAGE`

  The error number and error message of the most recent error that caused the I/O thread to stop. An error number of 0 and message of the empty string mean “no error.” If the `LAST_ERROR_MESSAGE` value is not empty, the error values also appear in the slave's error log.

  Issuing `RESET MASTER` or `RESET SLAVE` resets the values shown in these columns.

- `LAST_ERROR_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the most recent I/O error took place.

- `LAST_HEARTBEAT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the most recent heartbeat signal was received by a replication slave.

- `COUNT_RECEIVED_HEARTBEATS`

  The total number of heartbeat signals that a replication slave received since the last time it was restarted or reset, or a `CHANGE MASTER TO` statement was issued.

- `LAST_QUEUED_TRANSACTION`

  The global transaction ID (GTID) of the last transaction that was queued to the relay log.

- `LAST_QUEUED_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction queued in the relay log was committed on the original master.

- `LAST_QUEUED_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction queued in the relay log was committed on the immediate master.

- `LAST_QUEUED_TRANSACTION_START_QUEUE_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction was placed in the relay log queue by this I/O thread.

- `LAST_QUEUED_TRANSACTION_END_QUEUE_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction was queued to the relay log files.

- `QUEUEING_TRANSACTION`

  The global transaction ID (GTID) of the currently queueing transaction in the relay log.

- `QUEUEING_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the currently queueing transaction was committed on the original master.

- `QUEUEING_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the currently queueing transaction was committed on the immediate master.

- `QUEUEING_TRANSACTION_START_QUEUE_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the first event of the currently queueing transaction was written to the relay log by this I/O thread.

The `replication_connection_status` table has these indexes:

- Primary key on (`CHANNEL_NAME`)
- Index on (`THREAD_ID`)

`TRUNCATE TABLE` is not permitted for the `replication_connection_status` table.

The following table shows the correspondence between `replication_connection_status` columns and `SHOW SLAVE STATUS` columns.

`replication_connection_status` Column	**`SHOW SLAVE STATUS` Column**
`SOURCE_UUID`	`Master_UUID`

replication_connection_status Column	SHOW SLAVE STATUS Column
THREAD_ID	None
SERVICE_STATE	Slave_IO_Running
RECEIVED_TRANSACTION_SET	Retrieved_Gtid_Set
LAST_ERROR_NUMBER	Last_IO_Errno
LAST_ERROR_MESSAGE	Last_IO_Error
LAST_ERROR_TIMESTAMP	Last_IO_Error_Timestamp

25.11.11.3 The replication_applier_configuration Table

This table shows the configuration parameters that affect transactions applied by the slave server. Parameters stored in the table can be changed at runtime with the CHANGE MASTER TO statement, as indicated in the column descriptions.

The replication_applier_configuration table has these columns:

- CHANNEL_NAME

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

- DESIRED_DELAY

  The number of seconds that the slave must lag the master. (CHANGE MASTER TO option: MASTER_DELAY) See Section 18.3.11, “Delayed Replication” for more information.

The replication_applier_configuration table has these indexes:

- Primary key on (CHANNEL_NAME)

TRUNCATE TABLE is not permitted for the replication_applier_configuration table.

The following table shows the correspondence between replication_applier_configuration columns and SHOW SLAVE STATUS columns.

replication_applier_configuration Column	SHOW SLAVE STATUS Column
DESIRED_DELAY	SQL_Delay

25.11.11.4 The replication_applier_status Table

This table shows the current general transaction execution status on the slave server. The table provides information about general aspects of transaction applier status that are not specific to any thread involved. Thread-specific status information is available in the replication_applier_status_by_coordinator table (and replication_applier_status_by_worker if the slave is multi-threaded).

The replication_applier_status table has these columns:

- CHANNEL_NAME

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

- SERVICE_STATE

Shows ON when the replication channel's applier threads are active or idle, OFF means that the applier threads are not active.

- REMAINING_DELAY

  If the slave is waiting for DESIRED_DELAY seconds to pass since the master applied an transaction, this field contains the number of delay seconds remaining. At other times, this field is NULL. (The DESIRED_DELAY value is stored in the replication_applier_configuration table.) See Section 18.3.11, "Delayed Replication" for more information.

- COUNT_TRANSACTIONS_RETRIES

  Shows the number of retries that were made because the slave SQL thread failed to apply a transaction.

The replication_applier_status table has these indexes:

- Primary key on (CHANNEL_NAME)

TRUNCATE TABLE is not permitted for the replication_applier_status table.

The following table shows the correspondence between replication_applier_status columns and SHOW SLAVE STATUS columns.

replication_applier_status Column	SHOW SLAVE STATUS Column
SERVICE_STATE	None
REMAINING_DELAY	SQL_Remaining_Delay

25.11.11.5 The replication_applier_status_by_coordinator Table

For a multi-threaded slave, the slave uses multiple worker threads and a coordinator thread to manage them, and this table shows the status of the coordinator thread. For a single-threaded slave, this table is empty. For a multi-threaded slave, the replication_applier_status_by_worker table shows the status of the worker threads. This table provides information about the last transaction which was buffered by the coordinator thread to a worker's queue, as well as the transaction it is currently buffering. The start timestamp refers to when this thread read the first event of the transaction from the relay log to buffer it to a worker's queue, while the end timestamp refers to when the last event finished buffering to the worker's queue.

The replication_applier_status_by_coordinator table has these columns:

- CHANNEL_NAME

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, "Replication Channels" for more information.

- THREAD_ID

  The SQL/coordinator thread ID.

- SERVICE_STATE

  ON (thread exists and is active or idle) or OFF (thread no longer exists).

- LAST_ERROR_NUMBER, LAST_ERROR_MESSAGE

  The error number and error message of the most recent error that caused the SQL/coordinator thread to stop. An error number of 0 and message which is an empty string means "no error". If the LAST_ERROR_MESSAGE value is not empty, the error values also appear in the slave's error log.

Issuing `RESET MASTER` or `RESET SLAVE` resets the values shown in these columns.

All error codes and messages displayed in the `LAST_ERROR_NUMBER` and `LAST_ERROR_MESSAGE` columns correspond to error values listed in Section B.3, “Server Error Codes and Messages”.

- `LAST_ERROR_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the most recent SQL/coordinator error occurred.

- `LAST_PROCESSED_TRANSACTION`

  The global transaction ID (GTID) of the last transaction processed by this coordinator.

- `LAST_PROCESSED_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction processed by this coordinator was committed on the original master.

- `LAST_PROCESSED_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction processed by this coordinator was committed on the immediate master.

- `LAST_PROCESSED_TRANSACTION_START_BUFFER_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when this coordinator thread started writing the last transaction to the buffer of a worker thread.

- `LAST_PROCESSED_TRANSACTION_END_BUFFER_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction was written to the buffer of a worker thread by this coordinator thread.

- `PROCESSING_TRANSACTION`

  The global transaction ID (GTID) of the transaction that this coordinator thread is currently processing.

- `PROCESSING_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the currently processing transaction was committed on the original master.

- `PROCESSING_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the currently processing transaction was committed on the immediate master.

- `PROCESSING_TRANSACTION_START_BUFFER_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when this coordinator thread started writing the currently processing transaction to the buffer of a worker thread.

The `replication_applier_status_by_coordinator` table has these indexes:

- Primary key on (`CHANNEL_NAME`)
- Index on (`THREAD_ID`)

`TRUNCATE TABLE` is not permitted for the `replication_applier_status_by_coordinator` table.

The following table shows the correspondence between `replication_applier_status_by_coordinator` columns and `SHOW SLAVE STATUS` columns.

`replication_applier_status_by_coordinator` Column	`SHOW SLAVE STATUS` Column
`THREAD_ID`	None
`SERVICE_STATE`	`Slave_SQL_Running`
`LAST_ERROR_NUMBER`	`Last_SQL_Errno`
`LAST_ERROR_MESSAGE`	`Last_SQL_Error`
`LAST_ERROR_TIMESTAMP`	`Last_SQL_Error_Timestamp`

25.11.11.6 The replication_applier_status_by_worker Table

If the slave is not multi-threaded, this table shows the status of the applier thread. Otherwise, the slave uses multiple worker threads and a coordinator thread to manage them, and this table shows the status of the worker threads. For a multi-threaded slave, the `replication_applier_status_by_coordinator` table shows the status of the coordinator thread. This table provides details about the transactions applied, where the start timestamp refers to when the worker started applying the first event, and the end timestamp refers to when the last event of the transaction was applied.

The `replication_applier_status_by_worker` table has these columns:

- `CHANNEL_NAME`

  The replication channel which this row is displaying. There is always a default replication channel, and more replication channels can be added. See Section 18.2.3, “Replication Channels” for more information.

- `WORKER_ID`

  The worker identifier (same value as the `id` column in the `mysql.slave_worker_info` table). After `STOP SLAVE`, the `THREAD_ID` column becomes `NULL`, but the `WORKER_ID` value is preserved.

- `THREAD_ID`

  The worker thread ID.

- `SERVICE_STATE`

  `ON` (thread exists and is active or idle) or `OFF` (thread no longer exists).

- `LAST_ERROR_NUMBER`, `LAST_ERROR_MESSAGE`

  The error number and error message of the most recent error that caused the worker thread to stop. An error number of 0 and message of the empty string mean “no error”. If the `LAST_ERROR_MESSAGE` value is not empty, the error values also appear in the slave's error log.

  Issuing `RESET MASTER` or `RESET SLAVE` resets the values shown in these columns.

  All error codes and messages displayed in the `LAST_ERROR_NUMBER` and `LAST_ERROR_MESSAGE` columns correspond to error values listed in Section B.3, “Server Error Codes and Messages”.

- `LAST_ERROR_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the most recent worker error occurred.

- `LAST_APPLIED_TRANSACTION`

The global transaction ID (GTID) of the last transaction applied by this worker.

- `LAST_APPLIED_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction applied by this worker was committed on the original master.

- `LAST_APPLIED_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the last transaction applied by this worker was committed on the immediate master.

- `LAST_APPLIED_TRANSACTION_START_APPLY_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when this worker started applying the last applied transaction.

- `LAST_APPLIED_TRANSACTION_END_APPLY_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when this worker finished applying the last applied transaction.

- `APPLYING_TRANSACTION`

  The global transaction ID (GTID) of the transaction this worker is currently applying.

- `APPLYING_TRANSACTION_ORIGINAL_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the transaction this worker is currently applying was committed on the original master.

- `APPLYING_TRANSACTION_IMMEDIATE_COMMIT_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when the transaction this worker is currently applying was committed on the immediate master.

- `APPLYING_TRANSACTION_START_APPLY_TIMESTAMP`

  A timestamp in `'YYYY-MM-DD HH:MM:SS[.fraction]'` format that shows when this worker started applying the transaction that is currently being applied.

The `replication_applier_status_by_worker` table has these indexes:

- Primary key on (`CHANNEL_NAME`, `WORKER_ID`)
- Index on (`THREAD_ID`)

`TRUNCATE TABLE` is not permitted for the `replication_applier_status_by_worker` table.

The following table shows the correspondence between `replication_applier_status_by_worker` columns and `SHOW SLAVE STATUS` columns.

`replication_applier_status_by_worker` Column	**`SHOW SLAVE STATUS` Column**
`WORKER_ID`	None
`THREAD_ID`	None
`SERVICE_STATE`	None
`LAST_SEEN_TRANSACTION`	None
`LAST_ERROR_NUMBER`	`Last_SQL_Errno`

`replication_applier_status_by_worker` Column	`SHOW SLAVE STATUS` Column
`LAST_ERROR_MESSAGE`	`Last_SQL_Error`
`LAST_ERROR_TIMESTAMP`	`Last_SQL_Error_Timestamp`

25.11.11.7 The replication_applier_global_filters Table

This table shows the global replication filters configured on this slave. The `replication_applier_global_filters` table has the following columns:

- `FILTER_NAME`

  The type of replication filter that has been configured.

- `FILTER_RULE`

  The rules configured for the replication filter type using either `--replicate-*` command options or `CHANGE REPLICATION FILTER`.

- `CONFIGURED_BY`

  The method used to configure the replication filter, can be one of:

  - `CHANGE_REPLICATION_FILTER` configured by a global replication filter using a `CHANGE REPLICATION FILTER` statement.
  - `STARTUP_OPTIONS` configured by a global replication filter using a `--replicate-*` option.

- `ACTIVE_SINCE`

  Timestamp of when the replication filter was configured.

25.11.11.8 The replication_applier_filters Table

This table shows the replication channel specific filters configured on this slave. Each row provides information on a replication channel's configured type of filter. The `replication_applier_filters` table has the following columns:

- `CHANNEL_NAME`

  The name of replication channel with a replication filter configured.

- `FILTER_NAME`

  The type of replication filter that has been configured for this replication channel.

- `FILTER_RULE`

  The rules configured for the replication filter type using either `--replicate-*` command options or `CHANGE REPLICATION FILTER`.

- `CONFIGURED_BY`

  The method used to configure the replication filter, can be one of:

  - `CHANGE_REPLICATION_FILTER` configured by a global replication filter using a `CHANGE REPLICATION FILTER` statement.
  - `STARTUP_OPTIONS` configured by a global replication filter using a `--replicate-*` option.
  - `CHANGE_REPLICATION_FILTER_FOR_CHANNEL` configured by a channel specific replication filter using a `CHANGE REPLICATION FILTER FOR CHANNEL` statement.

  - `STARTUP_OPTIONS_FOR_CHANNEL` configured by a channel specific replication filter using a `--replicate-*` option.

- `ACTIVE_SINCE`

  Timestamp of when the replication filter was configured.

25.11.11.9 The replication_group_members Table

This table shows network and status information for replication group members.

The `replication_group_members` table has the following columns:

- `CHANNEL_NAME`

  Name of the group replication channel.

- `MEMBER_ID`

  Identifier for this member; same as the server UUID.

- `MEMBER_HOST`

  Network address of this member (host name or IP address).

- `MEMBER_PORT`

  Port on which the server is listening.

- `MEMBER_STATE`

  Current state of this member; can be any one of the following:

  - `OFFLINE`: The group replication plugin is installed but has not been started.
  - `RECOVERING`: The server has joined a group from which it is retrieving data.
  - `ONLINE`: The member is in a fully functioning state.

The `replication_group_members` table has these indexes:

- None

`TRUNCATE TABLE` is not permitted for the `replication_group_members` table.

25.11.11.10 The replication_group_member_stats Table

This table shows statistical information for replication group members.

The `replication_group_member_stats` table has the following columns:

- `CHANNEL_NAME`

  Name of the group replication channel

- `VIEW_ID`

  Current view identifier for this group.

- `MEMBER_ID`

  Identifier for this member; same as the server UUID.

- `COUNT_TRANSACTIONS_IN_QUEUE`

Number of transactions pending certification

- `COUNT_TRANSACTIONS_CHECKED`

  Number of transactions already certified by this member.

- `COUNT_CONFLICTS_DETECTED`

  Number of transactions that were negatively certified.

- `COUNT_TRANSACTIONS_VALIDATING`

  Number of transactions with which one can execute certification with them, but have not been garbage collected.

- `TRANSACTIONS_COMMITTED_ALL_MEMBERS`

  Set of stable group transactions.

- `LAST_CONFLICT_FREE_TRANSACTION`

  Latest transaction certified without conflicts.

The `replication_group_member_stats` table has these indexes:

- None

`TRUNCATE TABLE` is not permitted for the `replication_group_member_stats` table.

25.11.12 Performance Schema Lock Tables

The Performance Schema exposes lock information through these tables:

- `data_locks`: Data locks held and requested
- `data_lock_waits`: Relationships between data lock owners and data lock requestors blocked by those owners
- `metadata_locks`: Metadata locks held and requested
- `table_handles`: Table locks held and requested

The following sections describe these tables in more detail.

25.11.12.1 The data_locks Table

The `data_locks` table shows data locks held and requested. For information about which lock requests are blocked by which held locks, see Section 25.11.12.2, “The data_lock_waits Table”.

Example data lock information:

```
mysql> SELECT * FROM data_locks\G
*************************** 1. row ***************************
               ENGINE: INNODB
       ENGINE_LOCK_ID: 4140:74
ENGINE_TRANSACTION_ID: 4140
            THREAD_ID: 37
             EVENT_ID: 9
        OBJECT_SCHEMA: test
          OBJECT_NAME: t1
       PARTITION_NAME: NULL
    SUBPARTITION_NAME: NULL
           INDEX_NAME: NULL
OBJECT_INSTANCE_BEGIN: 140489308280888
            LOCK_TYPE: TABLE
```

```
           LOCK_MODE: IX
         LOCK_STATUS: GRANTED
           LOCK_DATA: NULL
*************************** 2. row ***************************
              ENGINE: INNODB
      ENGINE_LOCK_ID: 4140:66:5:1
ENGINE_TRANSACTION_ID: 4140
            THREAD_ID: 37
             EVENT_ID: 9
       OBJECT_SCHEMA: test
         OBJECT_NAME: t1
      PARTITION_NAME: NULL
   SUBPARTITION_NAME: NULL
          INDEX_NAME: GEN_CLUST_INDEX
OBJECT_INSTANCE_BEGIN: 140489320307736
            LOCK_TYPE: RECORD
            LOCK_MODE: X
         LOCK_STATUS: GRANTED
           LOCK_DATA: supremum pseudo-record
```

Unlike most Performance Schema data collection, there are no instruments for controlling whether data lock information is collected or system variables for controlling data lock table sizes. The Performance Schema collects information that is already available in the server, so there is no memory or CPU overhead to generate this information or need for parameters that control its collection.

Use the `data_locks` table to help diagnose performance problems that occur during times of heavy concurrent load. For `InnoDB`, see the discussion of this topic at Section 15.14.2, “InnoDB INFORMATION_SCHEMA Transaction and Locking Information”.

The `data_locks` table has these columns:

- `ENGINE`:

  The storage engine that holds or requested the lock.

- `ENGINE_LOCK_ID`:

  The ID of the lock held or requested by the storage engine. Tuples of (`ENGINE_LOCK_ID`, `ENGINE`) values are unique.

  Lock ID formats are internal and subject to change at any time. Applications should not rely on lock IDs having a particular format.

- `ENGINE_TRANSACTION_ID`:

  The storage engine internal ID of the transaction that requested the lock.

  For `InnoDB`, to obtain details about the transaction, join this column with the `TRX_ID` column of the `INFORMATION_SCHEMA INNODB_TRX` table.

- `THREAD_ID`:

  The thread ID of the that owns the lock. To obtain details about the thread, join this column with the `THREAD_ID` column of the Performance Schema `threads` table.

- `EVENT_ID`:

  The Performance Schema event that caused the lock. Tuples of (`THREAD_ID`, `EVENT_ID`) values implicitly identify a parent event in other Performance Schema tables:

  - The parent wait event in `events_waits_xxx` tables
  - The parent stage event in `events_stages_xxx` tables
  - The parent statement event in `events_statements_xxx` tables

  - The parent transaction event in `events_transactions_xxx` tables

  To obtain details about the parent event, join the `THREAD_ID` and `EVENT_ID` columns with the columns of like name in the appropriate parent event table.

- `OBJECT_SCHEMA`:

  The schema that contains the locked table.

- `OBJECT_NAME`:

  The name of the locked table.

- `PARTITION_NAME`:

  The name of the locked partition, if any; `NULL` otherwise.

- `SUBPARTITION_NAME`:

  The name of the locked subpartition, if any; `NULL` otherwise.

- `INDEX_NAME`:

  The name of the locked index, if any; `NULL` otherwise.

  In practice, `InnoDB` always creates an index (`GEN_CLUST_INDEX`), so `INDEX_NAME` is non-`NULL` for `InnoDB` tables.

- `OBJECT_INSTANCE_BEGIN`:

  The address in memory of the lock.

- `LOCK_TYPE`:

  The type of lock.

  The value is storage engine dependent. For `InnoDB`, permitted values are `RECORD` for a row-level lock, `TABLE` for a table-level lock.

- `LOCK_MODE`:

  How the lock is requested.

  The value is storage engine dependent. For `InnoDB`, permitted values are `S[,GAP]`, `X[,GAP]`, `IS[,GAP]`, `IX[,GAP]`, `AUTO_INC`, and `UNKNOWN`. Lock modes other than `AUTO_INC` and `UNKNOWN` indicate gap locks, if present. For information about `S`, `X`, `IS`, `IX`, and gap locks, refer to Section 15.5.1, “InnoDB Locking”.

- `LOCK_STATUS`:

  The status of the lock request.

  The value is storage engine dependent. For `InnoDB`, permitted values are `GRANTED` (lock is held) and `PENDING` (lock is being waited for).

- `LOCK_DATA`:

  The data associated with the lock, if any.

  The value is storage engine dependent. For `InnoDB`, values are primary key values of the locked record if `LOCK_TYPE` is `RECORD`, otherwise `NULL`. This column contains the values of the primary key columns in the locked row, formatted as a valid SQL string (ready to be copied to SQL statements).

If there is no primary key, `LOCK_DATA` is the unique `InnoDB` internal row ID number. If a gap lock is taken for key values or ranges above the largest value in the index, `LOCK_DATA` reports `supremum pseudo-record`. When the page containing the locked record is not in the buffer pool (in the case that it was paged out to disk while the lock was held), `InnoDB` does not fetch the page from disk, to avoid unnecessary disk operations. Instead, `LOCK_DATA` is set to `NULL`.

The `data_locks` table has these indexes:

- Primary key on (`ENGINE_LOCK_ID`, `ENGINE`)
- Index on (`ENGINE_TRANSACTION_ID`, `ENGINE`)
- Index on (`THREAD_ID`, `EVENT_ID`)
- Index on (`OBJECT_SCHEMA`, `OBJECT_NAME`, `PARTITION_NAME`, `SUBPARTITION_NAME`)

`TRUNCATE TABLE` is not permitted for the `data_locks` table.

25.11.12.2 The data_lock_waits Table

The `data_lock_waits` table implements a many-to-many relationship showing which data lock requests in the `data_locks` table are blocked by which held data locks in the `data_locks` table. Held locks in `data_locks` appear in `data_lock_waits` only if they block some lock request.

This information enables you to understand data lock dependencies between sessions. The table exposes not only which lock a session or transaction is waiting for, but which session or transaction currently holds that lock.

Example data lock wait information:

```
mysql> SELECT * FROM data_lock_waits\G
*************************** 1. row ***************************
                            ENGINE: INNODB
       REQUESTING_ENGINE_LOCK_ID: 4141:66:5:2
REQUESTING_ENGINE_TRANSACTION_ID: 4141
            REQUESTING_THREAD_ID: 38
             REQUESTING_EVENT_ID: 11
REQUESTING_OBJECT_INSTANCE_BEGIN: 140489320441368
         BLOCKING_ENGINE_LOCK_ID: 4140:66:5:2
  BLOCKING_ENGINE_TRANSACTION_ID: 4140
              BLOCKING_THREAD_ID: 37
               BLOCKING_EVENT_ID: 9
  BLOCKING_OBJECT_INSTANCE_BEGIN: 140489320307736
```

Unlike most Performance Schema data collection, there are no instruments for controlling whether data lock information is collected or system variables for controlling data lock table sizes. The Performance Schema collects information that is already available in the server, so there is no memory or CPU overhead to generate this information or need for parameters that control its collection.

Use the `data_lock_waits` table to help diagnose performance problems that occur during times of heavy concurrent load. For `InnoDB`, see the discussion of this topic at Section 15.14.2, “InnoDB INFORMATION_SCHEMA Transaction and Locking Information”.

Because the columns in the `data_lock_waits` table are similar to those in the `data_locks` table, the column descriptions here are abbreviated. For more detailed column descriptions, see Section 25.11.12.1, “The data_locks Table”.

The `data_lock_waits` table has these columns:

- `ENGINE`

  The storage engine that requested the lock.

- `REQUESTING_ENGINE_LOCK_ID`

The ID of the lock requested by the storage engine. To obtain details about the lock, join this column with the `ENGINE_LOCK_ID` column of the `data_locks` table.

- `REQUESTING_ENGINE_TRANSACTION_ID`

  The storage engine internal ID of the transaction that requested the lock.

- `REQUESTING_THREAD_ID`

  The thread ID of the session that requested the lock.

- `REQUESTING_EVENT_ID`

  The Performance Schema event that caused the lock request in the session that requested the lock.

- `REQUESTING_OBJECT_INSTANCE_BEGIN`

  The address in memory of the requested lock.

- `BLOCKING_ENGINE_LOCK_ID`

  The ID of the blocking lock. To obtain details about the lock, join this column with the `ENGINE_LOCK_ID` column of the `data_locks` table.

- `BLOCKING_ENGINE_TRANSACTION_ID`

  The storage engine internal ID of the transaction that holds the blocking lock.

- `BLOCKING_THREAD_ID`

  The thread ID of the session that holds the blocking lock.

- `BLOCKING_EVENT_ID`

  The Performance Schema event that caused the blocking lock in the session that holds it.

- `BLOCKING_OBJECT_INSTANCE_BEGIN`

  The address in memory of the blocking lock.

The `data_lock_waits` table has these indexes:

- Index on (`REQUESTING_ENGINE_LOCK_ID`, `ENGINE`)
- Index on (`BLOCKING_ENGINE_LOCK_ID`, `ENGINE`)
- Index on (`REQUESTING_ENGINE_TRANSACTION_ID`, `ENGINE`)
- Index on (`BLOCKING_ENGINE_TRANSACTION_ID`, `ENGINE`)
- Index on (`REQUESTING_THREAD_ID`, `REQUESTING_EVENT_ID`)
- Index on (`BLOCKING_THREAD_ID`, `BLOCKING_EVENT_ID`)

`TRUNCATE TABLE` is not permitted for the `data_lock_waits` table.

25.11.12.3 The metadata_locks Table

The Performance Schema exposes metadata lock information through the `metadata_locks` table:

- Locks that have been granted (shows which sessions own which current metadata locks)
- Locks that have been requested but not yet granted (shows which sessions are waiting for which metadata locks).

- Lock requests that have been killed by the deadlock detector or timed out and are waiting for the requesting session's lock request to be discarded

This information enables you to understand metadata lock dependencies between sessions. You can see not only which lock a session is waiting for, but which session currently holds that lock.

The `metadata_locks` table is read only and cannot be updated. It is autosized by default; to configure the table size, set the `performance_schema_max_metadata_locks` system variable at server startup.

Metadata lock instrumentation uses the `wait/lock/metadata/sql/mdl` instrument, which is enabled by default. (Prior to MySQL 8.0.2, the instrument is disabled by default.)

To control metadata lock instrumention state at server startup, use lines like these in your `my.cnf` file:

- Enable:

```
[mysqld]
performance-schema-instrument='wait/lock/metadata/sql/mdl=ON'
```

- Disable:

```
[mysqld]
performance-schema-instrument='wait/lock/metadata/sql/mdl=OFF'
```

To control metadata lock instrumention state at runtime, update the `setup_instruments` table:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
WHERE NAME = 'wait/lock/metadata/sql/mdl';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO', TIMED = 'NO'
WHERE NAME = 'wait/lock/metadata/sql/mdl';
```

The Performance Schema maintains `metadata_locks` table content as follows, using the `LOCK_STATUS` column to indicate the status of each lock:

- When a metadata lock is requested and obtained immediately, a row with a status of `GRANTED` is inserted.
- When a metadata lock is requested and not obtained immediately, a row with a status of `PENDING` is inserted.
- When a metadata lock previously requested is granted, its row status is updated to `GRANTED`.
- When a metadata lock is released, its row is deleted.
- When a pending lock request is canceled by the deadlock detector to break a deadlock (`ER_LOCK_DEADLOCK`), its row status is updated from `PENDING` to `VICTIM`.
- When a pending lock request times out (`ER_LOCK_WAIT_TIMEOUT`), its row status is updated from `PENDING` to `TIMEOUT`.
- When granted lock or pending lock request is killed, its row status is updated from `GRANTED` or `PENDING` to `KILLED`.
- The `VICTIM`, `TIMEOUT`, and `KILLED` status values are brief and signify that the lock row is about to be deleted.

- The PRE_ACQUIRE_NOTIFY and POST_RELEASE_NOTIFY status values are brief and signify that the metadata locking subsubsystem is notifying interested storage engines while entering lock acquisition or leaving lock release operations.

The metadata_locks table has these columns:

- OBJECT_TYPE

  The type of lock used in the metadata lock subsystem: The value is one of GLOBAL, SCHEMA, TABLE, FUNCTION, PROCEDURE, TRIGGER (currently unused), EVENT, COMMIT, USER LEVEL LOCK, TABLESPACE, or LOCKING SERVICE.

  A value of USER LEVEL LOCK indicates a lock acquired with GET_LOCK(). A value of LOCKING SERVICE indicates a lock acquired using the locking service described in Section 28.3.1, “The Locking Service”.

- OBJECT_SCHEMA

  The schema that contains the object.

- OBJECT_NAME

  The name of the instrumented object.

- OBJECT_INSTANCE_BEGIN

  The address in memory of the instrumented object.

- LOCK_TYPE

  The lock type from the metadata lock subsystem. The value is one of INTENTION_EXCLUSIVE, SHARED, SHARED_HIGH_PRIO, SHARED_READ, SHARED_WRITE, SHARED_UPGRADABLE, SHARED_NO_WRITE, SHARED_NO_READ_WRITE, or EXCLUSIVE.

- LOCK_DURATION

  The lock duration from the metadata lock subsystem. The value is one of STATEMENT, TRANSACTION, or EXPLICIT. The STATEMENT and TRANSACTION values signify locks that are released at statement or transaction end, respectively. The EXPLICIT value signifies locks that survive statement or transaction end and are released explicitly, such as global locks acquired with FLUSH TABLES WITH READ LOCK.

- LOCK_STATUS

  The lock status from the metadata lock subsystem. The value is one of PENDING, GRANTED, VICTIM, TIMEOUT, KILLED, PRE_ACQUIRE_NOTIFY, or POST_RELEASE_NOTIFY. The Performance Schema assigns these values as described previously.

- SOURCE

  The name of the source file containing the instrumented code that produced the event and the line number in the file at which the instrumentation occurs. This enables you to check the source to determine exactly what code is involved.

- OWNER_THREAD_ID

  The thread requesting a metadata lock.

- OWNER_EVENT_ID

  The event requesting a metadata lock.

The metadata_locks table has these indexes:

- Primary key on (OBJECT_INSTANCE_BEGIN)
- Index on (OBJECT_TYPE, OBJECT_SCHEMA, OBJECT_NAME)
- Index on (OWNER_THREAD_ID, OWNER_EVENT_ID)

TRUNCATE TABLE is not permitted for the metadata_locks table.

25.11.12.4 The table_handles Table

The Performance Schema exposes table lock information through the table_handles table to show the table locks currently in effect for each opened table handle. table_handles reports what is recorded by the table lock instrumentation. This information shows which table handles the server has open, how they are locked, and by which sessions.

The table_handles table is read only and cannot be updated. It is autosized by default; to configure the table size, set the performance_schema_max_table_handles system variable at server startup.

The table_handles table has these columns:

- OBJECT_TYPE

  The table opened by a table handle.

- OBJECT_SCHEMA

  The schema that contains the object.

- OBJECT_NAME

  The name of the instrumented object.

- OBJECT_INSTANCE_BEGIN

  The table handle address in memory.

- OWNER_THREAD_ID

  The thread owning the table handle.

- OWNER_EVENT_ID

  The event which caused the table handle to be opened.

- INTERNAL_LOCK

  The table lock used at the SQL level. The value is one of READ, READ WITH SHARED LOCKS, READ HIGH PRIORITY, READ NO INSERT, WRITE ALLOW WRITE, WRITE CONCURRENT INSERT, WRITE LOW PRIORITY, or WRITE. For information about these lock types, see the include/thr_lock.h source file.

- EXTERNAL_LOCK

  The table lock used at the storage engine level. The value is one of READ EXTERNAL or WRITE EXTERNAL.

The table_handles table has these indexes:

- Primary key on (OBJECT_INSTANCE_BEGIN)
- Index on (OBJECT_TYPE, OBJECT_SCHEMA, OBJECT_NAME)
- Index on (OWNER_THREAD_ID, OWNER_EVENT_ID)

`TRUNCATE TABLE` is not permitted for the `table_handles` table.

25.11.13 Performance Schema System Variable Tables

The MySQL server maintains many system variables that indicate how it is configured (see Section 5.1.5, “Server System Variables”). System variable information is available in these Performance Schema tables:

- `global_variables`: Global system variables. An application that wants only global values should use this table.
- `session_variables`: System variables for the current session. An application that wants all system variable values for its own session should use this table. It includes the session variables for its session, as well as the values of global variables that have no session counterpart.
- `variables_by_thread`: Session system variables for each active session. An application that wants to know the session variable values for specific sessions should use this table. It includes session variables only, identified by thread ID.
- `persisted_variables`: Provides a SQL interface to the `mysqld-auto.cnf` file that stores persisted global system variable settings. See Section 25.11.13.1, “Performance Schema persisted_variables Table”.
- `variables_info`: Shows, for each system variable, the source from which it was most recently set, and its range of values. See Section 25.11.13.2, “Performance Schema variables_info Table”.

The session variable tables (`session_variables`, `variables_by_thread`) contain information only for active sessions, not terminated sessions.

`TRUNCATE TABLE` is not supported for Performance Schema system variable tables.

The `global_variables` and `session_variables` tables have these columns:

- `VARIABLE_NAME`

  The system variable name.

- `VARIABLE_VALUE`

  The system variable value. For `global_variables`, this column contains the global value. For `session_variables`, this column contains the variable value in effect for the current session.

The `global_variables` and `session_variables` tables have these indexes:

- Primary key on (`VARIABLE_NAME`)

The `variables_by_thread` table has these columns:

- `THREAD_ID`

  The thread identifier of the session in which the system variable is defined.

- `VARIABLE_NAME`

  The system variable name.

- `VARIABLE_VALUE`

  The session variable value for the session named by the `THREAD_ID` column.

The `variables_by_thread` table has these indexes:

- Primary key on (`THREAD_ID`, `VARIABLE_NAME`)

The variables_by_thread table contains system variable information only about foreground threads. If not all threads are instrumented by the Performance Schema, this table will miss some rows. In this case, the Performance_schema_thread_instances_lost status variable will be greater than zero.

25.11.13.1 Performance Schema persisted_variables Table

The persisted_variables table provides an SQL interface to the mysqld-auto.cnf file that stores persisted global system variable settings, enabling the file contents to be inspected at runtime using SELECT statements. The table contains a row for each persisted system variable in the file. Variables not persisted do not appear in the table.

For information about persisted system variables, see Section 5.1.6, “Using System Variables”.

Suppose that mysqld-auto.cnf looks like this:

```
{ "mysql_server": {"autocommit": "ON" , "max_connections": "1000" } }
```

Then persisted_variables has these contents:

```
mysql> SELECT * FROM persisted_variables;
+-----------------+----------------+
| VARIABLE_NAME   | VARIABLE_VALUE |
+-----------------+----------------+
| autocommit      | ON             |
| max_connections | 1000           |
+-----------------+----------------+
```

persisted_variables has these columns:

- VARIABLE_NAME

  The variable name listed in mysqld-auto.cnf.

- VARIABLE_VALUE

  The value listed for the variable in mysqld-auto.cnf.

persisted_variables has these indexes:

- Primary key on (VARIABLE_NAME)

TRUNCATE TABLE is not permitted for the persisted_variables table.

25.11.13.2 Performance Schema variables_info Table

The variables_info table shows, for each system variable, the source from which it was most recently set, and its range of values.

variables_info has these columns:

- VARIABLE_NAME

  The variable name.

- VARIABLE_SOURCE

  The source from which the variable was most recently set:

  - COMMAND_LINE

    The variable was set on the command line.

  - `COMPILED`

    The variable has its compiled-in default value. `COMPILED` is the value used for variables not set any other way.

  - `DYNAMIC`

    The variable was set at runtime. This includes variables set within files specified using the `--init-file` option.

  - `EXPLICIT`

    The variable was set from an option file named with the `--defaults-file` option.

  - `EXTRA`

    The variable was set from an option file named with the `--defaults-extra-file` option.

  - `GLOBAL`

    The variable was set from a global option file. This includes option files not covered by `EXPLICIT`, `EXTRA`, `LOGIN`, `PERSISTED`, `SERVER`, or `USER`.

  - `LOGIN`

    The variable was set from a user-specific login path file (`~/.mylogin.cnf`).

  - `PERSISTED`

    The variable was set from a server-specific `mysqld-auto.cnf` option file. No row has this value if the server was started with `persisted_globals_load` disabled.

  - `SERVER`

    The variable was set from a server-specific `$MYSQL_HOME/my.cnf` option file. For details about how `MYSQL_HOME` is set, see Section 4.2.6, “Using Option Files”.

  - `USER`

    The variable was set from a user-specific `~/.my.cnf` option file.

- `VARIABLE_PATH`

  If the variable was set from an option file, `VARIABLE_PATH` is the path name of that file. Otherwise, the value is the empty string.

- `MIN_VALUE`, `MAX_VALUE`

  The minimum and maximum permitted values for the variable. Both are 0 for variables that have no such values (that is, variables that are not numeric).

- `SET_TIME`

  The time at which the variable was most recently set. The default is the time at which the server initialized global system variables during startup.

- `SET_USER`, `SET_HOST`

  The user name and host name of the client user that most recently set the variable.
  If a client connects as `user17` from host `host34.example.com` using the account `'user17'@'%.example.com`, `SET_USER` and `SET_HOST` will be `user17` and `host34.example.com`, respectively. For proxy user connections, these values correspond to the

external (proxy) user, not the proxied user against which privilege checking is performed. The default for each column is the empty string, indicating that the variable has not been set since server startup.

The `variables_info` table has these indexes:

- None

`TRUNCATE TABLE` is not permitted for the `variables_info` table.

If a variable with a `VARIABLE_SOURCE` value other than `DYNAMIC` is set at runtime, `VARIABLE_SOURCE` becomes `DYNAMIC` and `VARIABLE_PATH` becomes the empty string.

A system variable that has only a session value (such as `debug_sync`) cannot be set at startup or persisted. For session-only system variables, `VARIABLE_SOURCE` can be only `COMPILED` or `DYNAMIC`.

If a system variable has an unexpected `VARIABLE_SOURCE` value, consider your server startup method. For example, `mysqld_safe` reads option files and passes certain options it finds there as part of the command line that it uses to start `mysqld`. Consequently, some system variables that you set in option files might display in `variables_info` as `COMMAND_LINE`, rather than as `GLOBAL` or `SERVER` as you might otherwise expect.

Some sample queries that use the `variables_info` table, with representive output:

- Display variables set on the command line:

```
mysql> SELECT VARIABLE_NAME FROM variables_info
       WHERE VARIABLE_SOURCE = 'COMMAND_LINE'
       ORDER BY VARIABLE_NAME;
+---------------+
| VARIABLE_NAME |
+---------------+
| basedir       |
| datadir       |
| log_error     |
| pid_file      |
| plugin_dir    |
| port          |
+---------------+
```

- Display variables set from persistent storage:

```
mysql> SELECT VARIABLE_NAME FROM variables_info
       WHERE VARIABLE_SOURCE = 'PERSISTED'
       ORDER BY VARIABLE_NAME;
+--------------------------+
| VARIABLE_NAME            |
+--------------------------+
| event_scheduler          |
| expire_logs_days         |
| max_connections          |
| validate_password_policy |
+--------------------------+
```

- Join `variables_info` with the `global_variables` table to display the current values of persisted variables, together with their range of values:

```
mysql> SELECT
         VI.VARIABLE_NAME, GV.VARIABLE_VALUE,
         VI.MIN_VALUE,VI.MAX_VALUE
       FROM variables_info AS VI
         INNER JOIN global_variables AS GV
         USING(VARIABLE_NAME)
       WHERE VI.VARIABLE_SOURCE = 'PERSISTED'
       ORDER BY VARIABLE_NAME;
+--------------------------+----------------+-----------+-----------+
```

```
| VARIABLE_NAME            | VARIABLE_VALUE | MIN_VALUE | MAX_VALUE |
+--------------------------+----------------+-----------+-----------+
| event_scheduler          | ON             | 0         | 0         |
| expire_logs_days         | 7              | 0         | 99        |
| max_connections          | 200            | 1         | 100000    |
| validate_password_policy | STRONG         | 0         | 0         |
+--------------------------+----------------+-----------+-----------+
```

25.11.14 Performance Schema Status Variable Tables

The MySQL server maintains many status variables that provide information about its operation (see Section 5.1.7, “Server Status Variables”). Status variable information is available in these Performance Schema tables:

- `global_status`: Global status variables. An application that wants only global values should use this table.
- `session_status`: Status variables for the current session. An application that wants all status variable values for its own session should use this table. It includes the session variables for its session, as well as the values of global variables that have no session counterpart.
- `status_by_thread`: Session status variables for each active session. An application that wants to know the session variable values for specific sessions should use this table. It includes session variables only, identified by thread ID.

There are also summary tables that provide status variable information aggregated by account, host name, and user name. See Section 25.11.15.12, “Status Variable Summary Tables”.

The session variable tables (`session_status`, `status_by_thread`) contain information only for active sessions, not terminated sessions.

The Performance Schema collects statistics for global status variables only for threads for which the `INSTRUMENTED` value is `YES` in the `threads` table. Statistics for session status variables are always collected, regardless of the `INSTRUMENTED` value.

The Performance Schema does not collect statistics for `Com_xxx` status variables in the status variable tables. To obtain global and per-session statement execution counts, use the `events_statements_summary_global_by_event_name` and `events_statements_summary_by_thread_by_event_name` tables, respectively. For example:

```
SELECT EVENT_NAME, COUNT_STAR
FROM events_statements_summary_global_by_event_name
WHERE EVENT_NAME LIKE 'statement/sql/%';
```

The `global_status` and `session_status` tables have these columns:

- `VARIABLE_NAME`

  The status variable name.

- `VARIABLE_VALUE`

  The status variable value. For `global_status`, this column contains the global value. For `session_status`, this column contains the variable value for the current session.

The `global_status` and `session_status` tables have these indexes:

- Primary key on (`VARIABLE_NAME`)

The `status_by_thread` table contains the status of each active thread. It has these columns:

- `THREAD_ID`

  The thread identifier of the session in which the status variable is defined.

- `VARIABLE_NAME`

  The status variable name.

- `VARIABLE_VALUE`

  The session variable value for the session named by the `THREAD_ID` column.

The `status_by_thread` table has these indexes:

- Primary key on (`THREAD_ID`, `VARIABLE_NAME`)

The `status_by_thread` table contains status variable information only about foreground threads. If the `performance_schema_max_thread_instances` system variable is not autoscaled (set to −1) and the maximum permitted number of instrumented thread objects is not greater than the number of background threads, the table will be empty.

The Performance Schema supports `TRUNCATE TABLE` for status variable tables as follows:

- `global_status`: Resets thread, account, host, and user status. Resets global status variables except those that the server never resets.
- `session_status`: Not supported.
- `status_by_thread`: Aggregates status for all threads to the global status and account status, then resets thread status. If account statistics are not collected, the session status is added to host and user status, if host and user status are collected.

  Account, host, and user statistics are not collected if the `performance_schema_accounts_size`, `performance_schema_hosts_size`, and `performance_schema_users_size` system variables, respectively, are set to 0.

`FLUSH STATUS` adds the session status from all active sessions to the global status variables, resets the status of all active sessions, and resets account, host, and user status values aggregated from disconnected sessions.

25.11.15 Performance Schema Summary Tables

Summary tables provide aggregated information for terminated events over time. The tables in this group summarize event data in different ways.

Wait Event Summaries

- `events_waits_summary_by_account_by_event_name`: Wait events per account and event name
- `events_waits_summary_by_host_by_event_name`: Wait events per host name and event name
- `events_waits_summary_by_instance`: Wait events per instance
- `events_waits_summary_by_thread_by_event_name`: Wait events per thread and event name
- `events_waits_summary_by_user_by_event_name`: Wait events per user name and event name
- `events_waits_summary_global_by_event_name`: Wait events per event name

Stage Summaries

- `events_stages_summary_by_account_by_event_name`: Stage events per account and event name

- `events_stages_summary_by_host_by_event_name`: Stage events per host name and event name
- `events_stages_summary_by_thread_by_event_name`: Stage waits per thread and event name
- `events_stages_summary_by_user_by_event_name`: Stage events per user name and event name
- `events_stages_summary_global_by_event_name`: Stage waits per event name

Statement Summaries

- `events_statements_histogram_by_digest`: Statement histograms per schema and digest value.
- `events_statements_histogram_global`: Statement histogram summarized globally.
- `events_statements_summary_by_account_by_event_name`: Statement events per account and event name
- `events_statements_summary_by_digest`: Statement events per schema and digest value
- `events_statements_summary_by_host_by_event_name`: Statement events per host name and event name
- `events_statements_summary_by_program`: Statement events per stored program (stored procedures and functions, triggers, and events)
- `events_statements_summary_by_thread_by_event_name`: Statement events per thread and event name
- `events_statements_summary_by_user_by_event_name`: Statement events per user name and event name
- `events_statements_summary_global_by_event_name`: Statement events per event name
- `prepared_statements_instances`: Prepared statement instances and statistics

Transaction Summaries

- `events_transactions_summary_by_account_by_event_name`: Transaction events per account and event name
- `events_transactions_summary_by_host_by_event_name`: Transaction events per host name and event name
- `events_transactions_summary_by_thread_by_event_name`: Transaction events per thread and event name
- `events_transactions_summary_by_user_by_event_name`: Transaction events per user name and event name
- `events_transactions_summary_global_by_event_name`: Transaction events per event name

Object Wait Summaries

- `objects_summary_global_by_type`: Object summaries

File I/O Summaries

- `file_summary_by_event_name`: File events per event name

- `file_summary_by_instance`: File events per file instance

Table I/O and Lock Wait Summaries

- `table_io_waits_summary_by_index_usage`: Table I/O waits per index
- `table_io_waits_summary_by_table`: Table I/O waits per table
- `table_lock_waits_summary_by_table`: Table lock waits per table

Socket Summaries

- `socket_summary_by_instance`: Socket waits and I/O per instance
- `socket_summary_by_event_name`: Socket waits and I/O per event name

Memory Summaries

- `memory_summary_by_account_by_event_name`: Memory operations per account and event name
- `memory_summary_by_host_by_event_name`: Memory operations per host and event name
- `memory_summary_by_thread_by_event_name`: Memory operations per thread and event name
- `memory_summary_by_user_by_event_name`: Memory operations per user and event name
- `memory_summary_global_by_event_name`: Memory operations globally per event name

Error Summaries

- `events_errors_summary_by_account_by_error`: Errors per error code and account
- `events_errors_summary_by_host_by_error`: Errors per error code and host
- `events_errors_summary_by_thread_by_error`: Errors per error code and thread
- `events_errors_summary_by_user_by_error`: Errors per error code and user
- `events_errors_summary_global_by_error`: Errors per error code

Status Variable Summaries

- `status_by_account`: Status variables per account
- `status_by_host`: Status variables per host name
- `status_by_user`: Status variables per user name

Each summary table has grouping columns that determine how to group the data to be aggregated, and summary columns that contain the aggregated values. Tables that summarize events in similar ways often have similar sets of summary columns and differ only in the grouping columns used to determine how events are aggregated.

Summary tables can be truncated with `TRUNCATE TABLE`. Generally, the effect is to reset the summary columns to 0 or `NULL`, not to remove rows. This enables you to clear collected values and restart aggregation. That might be useful, for example, after you have made a runtime configuration change. Exceptions to this truncation behavior are noted in individual summary table sections.

25.11.15.1 Wait Event Summary Tables

The Performance Schema maintains tables for collecting current and recent wait events, and aggregates that information in summary tables. Section 25.11.4, “Performance Schema Wait Event

Tables" describes the events on which wait summaries are based. See that discussion for information about the content of wait events, the current and recent wait event tables, and how to control wait event collection, which is disabled by default.

Example wait event summary information:

```
mysql> SELECT * FROM events_waits_summary_global_by_event_name\G
...
*************************** 6. row ***************************
    EVENT_NAME: wait/synch/mutex/sql/BINARY_LOG::LOCK_index
    COUNT_STAR: 8
SUM_TIMER_WAIT: 2119302
MIN_TIMER_WAIT: 196092
AVG_TIMER_WAIT: 264912
MAX_TIMER_WAIT: 569421
...
*************************** 9. row ***************************
    EVENT_NAME: wait/synch/mutex/sql/hash_filo::lock
    COUNT_STAR: 69
SUM_TIMER_WAIT: 16848828
MIN_TIMER_WAIT: 0
AVG_TIMER_WAIT: 244185
MAX_TIMER_WAIT: 735345
...
```

Each wait event summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `events_waits_summary_by_account_by_event_name` has `EVENT_NAME`, `USER`, and `HOST` columns. Each row summarizes events for a given account (user and host combination) and event name.
- `events_waits_summary_by_host_by_event_name` has `EVENT_NAME` and `HOST` columns. Each row summarizes events for a given host and event name.
- `events_waits_summary_by_instance` has `EVENT_NAME` and `OBJECT_INSTANCE_BEGIN` columns. Each row summarizes events for a given event name and object. If an instrument is used to create multiple instances, each instance has a unique `OBJECT_INSTANCE_BEGIN` value and is summarized separately in this table.
- `events_waits_summary_by_thread_by_event_name` has `THREAD_ID` and `EVENT_NAME` columns. Each row summarizes events for a given thread and event name.
- `events_waits_summary_by_user_by_event_name` has `EVENT_NAME` and `USER` columns. Each row summarizes events for a given user and event name.
- `events_waits_summary_global_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name. An instrument might be used to create multiple instances of the instrumented object. For example, if there is an instrument for a mutex that is created for each connection, there are as many instances as there are connections. The summary row for the instrument summarizes over all these instances.

Each wait event summary table has these summary columns containing aggregated values:

- `COUNT_STAR`

  The number of summarized events. This value includes all events, whether timed or nontimed.

- `SUM_TIMER_WAIT`

  The total wait time of the summarized timed events. This value is calculated only for timed events because nontimed events have a wait time of `NULL`. The same is true for the other *`xxx`*`_TIMER_WAIT` values.

- `MIN_TIMER_WAIT`

  The minimum wait time of the summarized timed events.

- `AVG_TIMER_WAIT`

  The average wait time of the summarized timed events.

- `MAX_TIMER_WAIT`

  The maximum wait time of the summarized timed events.

The wait event summary tables have these indexes:

- `events_waits_summary_by_account_by_event_name`:
  - Primary key on (`USER`, `HOST`, `EVENT_NAME`)
- `events_waits_summary_by_host_by_event_name`:
  - Primary key on (`HOST`, `EVENT_NAME`)
- `events_waits_summary_by_instance`:
  - Primary key on (`OBJECT_INSTANCE_BEGIN`)
  - Index on (`EVENT_NAME`)
- `events_waits_summary_by_thread_by_event_name`:
  - Primary key on (`THREAD_ID`, `EVENT_NAME`)
- `events_waits_summary_by_user_by_event_name`:
  - Primary key on (`USER`, `EVENT_NAME`)
- `events_waits_summary_global_by_event_name`:
  - Primary key on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for wait summary tables. It has these effects:

- For summary tables not aggregated by account, host, or user, truncation resets the summary columns to zero rather than removing rows.
- For summary tables aggregated by account, host, or user, truncation removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero for the remaining rows.

In addition, each wait summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of `events_waits_summary_global_by_event_name`. For details, see Section 25.11.8, “Performance Schema Connection Tables”.

25.11.15.2 Stage Summary Tables

The Performance Schema maintains tables for collecting current and recent stage events, and aggregates that information in summary tables. Section 25.11.5, “Performance Schema Stage Event Tables” describes the events on which stage summaries are based. See that discussion for information about the content of stage events, the current and recent stage event tables, and how to control stage event collection, which is disabled by default.

Example stage event summary information:

```
mysql> SELECT * FROM events_stages_summary_global_by_event_name\G
```

```
...
*************************** 5. row ***************************
    EVENT_NAME: stage/sql/checking permissions
    COUNT_STAR: 57
SUM_TIMER_WAIT: 26501888880
MIN_TIMER_WAIT: 7317456
AVG_TIMER_WAIT: 464945295
MAX_TIMER_WAIT: 12858936792
...
*************************** 9. row ***************************
    EVENT_NAME: stage/sql/closing tables
    COUNT_STAR: 37
SUM_TIMER_WAIT: 662606568
MIN_TIMER_WAIT: 1593864
AVG_TIMER_WAIT: 17907891
MAX_TIMER_WAIT: 437977248
...
```

Each stage summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `events_stages_summary_by_account_by_event_name` has `EVENT_NAME`, `USER`, and `HOST` columns. Each row summarizes events for a given account (user and host combination) and event name.
- `events_stages_summary_by_host_by_event_name` has `EVENT_NAME` and `HOST` columns. Each row summarizes events for a given host and event name.
- `events_stages_summary_by_thread_by_event_name` has `THREAD_ID` and `EVENT_NAME` columns. Each row summarizes events for a given thread and event name.
- `events_stages_summary_by_user_by_event_name` has `EVENT_NAME` and `USER` columns. Each row summarizes events for a given user and event name.
- `events_stages_summary_global_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.

Each stage summary table has these summary columns containing aggregated values: `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, and `MAX_TIMER_WAIT`. These columns are analogous to the columns of the same names in the wait event summary tables (see Section 25.11.15.1, “Wait Event Summary Tables”), except that the stage summary tables aggregate events from `events_stages_current` rather than `events_waits_current`.

The stage summary tables have these indexes:

- `events_stages_summary_by_account_by_event_name`:
  - Primary key on (`USER`, `HOST`, `EVENT_NAME`)
- `events_stages_summary_by_host_by_event_name`:
  - Primary key on (`HOST`, `EVENT_NAME`)
- `events_stages_summary_by_thread_by_event_name`:
  - Primary key on (`THREAD_ID`, `EVENT_NAME`)
- `events_stages_summary_by_user_by_event_name`:
  - Primary key on (`USER`, `EVENT_NAME`)
- `events_stages_summary_global_by_event_name`:
  - Primary key on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for stage summary tables. It has these effects:

- For summary tables not aggregated by account, host, or user, truncation resets the summary columns to zero rather than removing rows.
- For summary tables aggregated by account, host, or user, truncation removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero for the remaining rows.

In addition, each stage summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of `events_stages_summary_global_by_event_name`. For details, see Section 25.11.8, “Performance Schema Connection Tables”.

25.11.15.3 Statement Summary Tables

The Performance Schema maintains tables for collecting current and recent statement events, and aggregates that information in summary tables. Section 25.11.6, “Performance Schema Statement Event Tables” describes the events on which statement summaries are based. See that discussion for information about the content of statement events, the current and recent statement event tables, and how to control statement event collection, which is partially disabled by default.

Example statement event summary information:

```
mysql> SELECT * FROM events_statements_summary_global_by_event_name\G
*************************** 1. row ***************************
                EVENT_NAME: statement/sql/select
                COUNT_STAR: 25
            SUM_TIMER_WAIT: 1535983999000
            MIN_TIMER_WAIT: 209823000
            AVG_TIMER_WAIT: 61439359000
            MAX_TIMER_WAIT: 1363397650000
             SUM_LOCK_TIME: 20186000000
                SUM_ERRORS: 0
              SUM_WARNINGS: 0
         SUM_ROWS_AFFECTED: 0
             SUM_ROWS_SENT: 388
         SUM_ROWS_EXAMINED: 370
SUM_CREATED_TMP_DISK_TABLES: 0
     SUM_CREATED_TMP_TABLES: 0
       SUM_SELECT_FULL_JOIN: 0
 SUM_SELECT_FULL_RANGE_JOIN: 0
           SUM_SELECT_RANGE: 0
     SUM_SELECT_RANGE_CHECK: 0
            SUM_SELECT_SCAN: 6
      SUM_SORT_MERGE_PASSES: 0
             SUM_SORT_RANGE: 0
              SUM_SORT_ROWS: 0
              SUM_SORT_SCAN: 0
          SUM_NO_INDEX_USED: 6
     SUM_NO_GOOD_INDEX_USED: 0
...
```

Each statement summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `events_statements_summary_by_account_by_event_name` has `EVENT_NAME`, `USER`, and `HOST` columns. Each row summarizes events for a given account (user and host combination) and event name.
- `events_statements_summary_by_digest` has `SCHEMA_NAME` and `DIGEST` columns. Each row summarizes events per schema and digest value. (The `DIGEST_TEXT` column contains the corresponding normalized statement digest text, but is neither a grouping nor summary column.)

  The maximum number of rows in the table is autosized at server startup. To set this maximum explicitly, set the `performance_schema_digests_size` system variable at server startup.

- `events_statements_summary_by_host_by_event_name` has `EVENT_NAME` and `HOST` columns. Each row summarizes events for a given host and event name.
- `events_statements_summary_by_program` has `OBJECT_TYPE`, `OBJECT_SCHEMA`, and `OBJECT_NAME` columns. Each row summarizes events for a given stored program (stored procedure or function, trigger, or event).
- `events_statements_summary_by_thread_by_event_name` has `THREAD_ID` and `EVENT_NAME` columns. Each row summarizes events for a given thread and event name.
- `events_statements_summary_by_user_by_event_name` has `EVENT_NAME` and `USER` columns. Each row summarizes events for a given user and event name.
- `events_statements_summary_global_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.
- `prepared_statements_instances` has an `OBJECT_INSTANCE_BEGIN` column. Each row summarizes events for a given prepared statement.

Each statement summary table has these summary columns containing aggregated values (with exceptions as noted):

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns are analogous to the columns of the same names in the wait event summary tables (see Section 25.11.15.1, “Wait Event Summary Tables”), except that the statement summary tables aggregate events from `events_statements_current` rather than `events_waits_current`.

  The `prepared_statements_instances` table does not have these columns.

- `SUM_`*`xxx`*

  The aggregate of the corresponding *`xxx`* column in the `events_statements_current` table. For example, the `SUM_LOCK_TIME` and `SUM_ERRORS` columns in statement summary tables are the aggregates of the `LOCK_TIME` and `ERRORS` columns in `events_statements_current` table.

The `events_statements_summary_by_digest` table has these additional summary columns:

- `FIRST_SEEN_TIMESTAMP`, `LAST_SEEN_TIMESTAMP`

  The times at which a statement with the given digest value were first seen and most recently seen.

- `QUANTILE_95`: The 95th percentile of the statement latency, in picoseconds. This percentile is a high estimate, computed from the histogram data collected. In other words, for a given digest, 95% of the statements measured have a latency lower than `QUANTILE_95`.

  For access to the histogram data, use the tables described in Section 25.11.15.4, “Statement Histogram Summary Tables”.

- `QUANTILE_99`: Similar to `QUANTILE_95`, but for the 99th percentile.
- `QUANTILE_999`: Similar to `QUANTILE_95`, but for the 99.9th percentile.

The `events_statements_summary_by_program` table has these additional summary columns:

- `COUNT_STATEMENTS`, `SUM_STATEMENTS_WAIT`, `MIN_STATEMENTS_WAIT`, `AVG_STATEMENTS_WAIT`, `MAX_STATEMENTS_WAIT`

  Statistics about nested statements invoked during stored program execution.

The `prepared_statements_instances` table has these additional summary columns:

- `COUNT_EXECUTE`, `SUM_TIMER_EXECUTE`, `MIN_TIMER_EXECUTE`, `AVG_TIMER_EXECUTE`, `MAX_TIMER_EXECUTE`

Aggregated statistics for executions of the prepared statement.

The statement summary tables have these indexes:

- `events_transactions_summary_by_account_by_event_name`:
  - Primary key on (`USER`, `HOST`, `EVENT_NAME`)
- `events_statements_summary_by_digest`:
  - Primary key on (`SCHEMA_NAME`, `DIGEST`)
- `events_transactions_summary_by_host_by_event_name`:
  - Primary key on (`HOST`, `EVENT_NAME`)
- `events_statements_summary_by_program`:
  - Primary key on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`)
- `events_statements_summary_by_thread_by_event_name`:
  - Primary key on (`THREAD_ID`, `EVENT_NAME`)
- `events_transactions_summary_by_user_by_event_name`:
  - Primary key on (`USER`, `EVENT_NAME`)
- `events_statements_summary_global_by_event_name`:
  - Primary key on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for statement summary tables. It has these effects:

- For `events_statements_summary_by_digest`, it removes the rows.
- For other summary tables not aggregated by account, host, or user, truncation resets the summary columns to zero rather than removing rows.
- For other summary tables aggregated by account, host, or user, truncation removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero for the remaining rows.

In addition, each statement summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of `events_statements_summary_global_by_event_name`. For details, see Section 25.11.8, “Performance Schema Connection Tables”.

In addition, truncating `events_statements_summary_by_digest` implicitly truncates `events_statements_histogram_by_digest`, and truncating `events_statements_summary_global_by_event_name` implicitly truncates `events_statements_histogram_global`.

Statement Digest Aggregation Rules

If the `statements_digest` consumer is enabled, aggregation into `events_statements_summary_by_digest` occurs as follows when a statement completes. Aggregation is based on the `DIGEST` value computed for the statement.

- If a `events_statements_summary_by_digest` row already exists with the digest value for the statement that just completed, statistics for the statement are aggregated to that row. The `LAST_SEEN` column is updated to the current time.

- If no row has the digest value for the statement that just completed, and the table is not full, a new row is created for the statement. The `FIRST_SEEN` and `LAST_SEEN` columns are initialized with the current time.
- If no row has the statement digest value for the statement that just completed, and the table is full, the statistics for the statement that just completed are added to a special "catch-all" row with `DIGEST = NULL`, which is created if necessary. If the row is created, the `FIRST_SEEN` and `LAST_SEEN` columns are initialized with the current time. Otherwise, the `LAST_SEEN` column is updated with the current time.

The row with `DIGEST = NULL` is maintained because Performance Schema tables have a maximum size due to memory constraints. The `DIGEST = NULL` row permits digests that do not match other rows to be counted even if the summary table is full, using a common "other" bucket. This row helps you estimate whether the digest summary is representative:

- A `DIGEST = NULL` row that has a `COUNT_STAR` value that represents 5% of all digests shows that the digest summary table is very representative; the other rows cover 95% of the statements seen.
- A `DIGEST = NULL` row that has a `COUNT_STAR` value that represents 50% of all digests shows that the digest summary table is not very representative; the other rows cover only half the statements seen. Most likely the DBA should increase the maximum table size so that more of the rows counted in the `DIGEST = NULL` row would be counted using more specific rows instead. To do this, set the `performance_schema_digests_size` system variable to a larger value at server startup. The default size is 200.

Stored Program Instrumentation Behavior

For stored program types for which instrumentation is enabled in the `setup_objects` table, `events_statements_summary_by_program` maintains statistics for stored programs as follows:

- A row is added for an object when it is first used in the server.
- The row for an object is removed when the object is dropped.
- Statistics are aggregated in the row for an object as it executes.

See also Section 25.4.3, "Event Pre-Filtering".

25.11.15.4 Statement Histogram Summary Tables

The Performance Schema maintains statement event summary tables that contain information about minimum, maximum, and average statement latency (see Section 25.11.15.3, "Statement Summary Tables"). Those tables permit high-level assessment of system performance. To permit assessment at a more fine-grained level, the Performance Schema also collects histogram data for statement latencies. These histograms provide additional insight into latency distributions.

Section 25.11.6, "Performance Schema Statement Event Tables" describes the events on which statement summaries are based. See that discussion for information about the content of statement events, the current and recent statement event tables, and how to control statement event collection, which is partially disabled by default.

Example statement histogram information:

```
mysql> SELECT * FROM events_statements_histogram_by_digest
       WHERE SCHEMA_NAME = 'mydb' AND DIGEST = 'bb3f69453119b2d7b3ae40673a9d4c7c'
       AND COUNT_BUCKET > 0 ORDER BY BUCKET_NUMBER\G
*************************** 1. row ***************************
           SCHEMA_NAME: mydb
                DIGEST: bb3f69453119b2d7b3ae40673a9d4c7c
         BUCKET_NUMBER: 42
      BUCKET_TIMER_LOW: 66069344
     BUCKET_TIMER_HIGH: 69183097
```

```
          COUNT_BUCKET: 1
COUNT_BUCKET_AND_LOWER: 1
       BUCKET_QUANTILE: 0.058824
*************************** 2. row ***************************
           SCHEMA_NAME: mydb
                DIGEST: bb3f69453119b2d7b3ae40673a9d4c7c
         BUCKET_NUMBER: 43
      BUCKET_TIMER_LOW: 69183097
     BUCKET_TIMER_HIGH: 72443596
          COUNT_BUCKET: 1
COUNT_BUCKET_AND_LOWER: 2
       BUCKET_QUANTILE: 0.117647
*************************** 3. row ***************************
           SCHEMA_NAME: mydb
                DIGEST: bb3f69453119b2d7b3ae40673a9d4c7c
         BUCKET_NUMBER: 44
      BUCKET_TIMER_LOW: 72443596
     BUCKET_TIMER_HIGH: 75857757
          COUNT_BUCKET: 2
COUNT_BUCKET_AND_LOWER: 4
       BUCKET_QUANTILE: 0.235294
*************************** 4. row ***************************
           SCHEMA_NAME: mydb
                DIGEST: bb3f69453119b2d7b3ae40673a9d4c7c
         BUCKET_NUMBER: 45
      BUCKET_TIMER_LOW: 75857757
     BUCKET_TIMER_HIGH: 79432823
          COUNT_BUCKET: 6
COUNT_BUCKET_AND_LOWER: 10
       BUCKET_QUANTILE: 0.625000
...
```

For example, in row 3, these values indicate that 23.52% of queries run in under 75.86 microseconds:

```
BUCKET_TIMER_HIGH: 75857757
  BUCKET_QUANTILE: 0.235294
```

In row 4, these values indicate that 62.50% of queries run in under 79.44 microseconds:

```
BUCKET_TIMER_HIGH: 79432823
  BUCKET_QUANTILE: 0.625000
```

Each statement histogram summary table has one or more grouping columns to indicate how the table aggregates events:

- `events_statements_histogram_by_digest` has `SCHEMA_NAME`, `DIGEST`, and `BUCKET_NUMBER` columns:
  - The `SCHEMA_NAME` and `DIGEST` columns identify a statement digest row in the `events_statements_summary_by_digest` table.
  - The `events_statements_histogram_by_digest` rows with the same `SCHEMA_NAME` and `DIGEST` values comprise the histogram for that schema/digest combination.
  - Within a given histogram, the `BUCKET_NUMBER` column indicates the bucket number.
- `events_statements_histogram_global` has a `BUCKET_NUMBER` column. This table summarizes latencies globally across schema name and digest values, using a single histogram. The `BUCKET_NUMBER` column indicates the bucket number within this global histogram.

A histogram consists of *`N`* buckets, where each row represents one bucket, with the bucket number indicated by the `BUCKET_NUMBER` column. Bucket numbers begin with 0.

Each statement histogram summary table has these summary columns containing aggregated values:

- `BUCKET_TIMER_LOW`, `BUCKET_TIMER_HIGH`

A bucket counts statements that have a latency, in picoseconds, measured between `BUCKET_TIMER_LOW` and `BUCKET_TIMER_HIGH`:

- The value of `BUCKET_TIMER_LOW` for the first bucket (`BUCKET_NUMBER` = 0) is 0.
- The value of `BUCKET_TIMER_LOW` for a bucket (`BUCKET_NUMBER` = *k*) is the same as `BUCKET_TIMER_HIGH` for the previous bucket (`BUCKET_NUMBER` = *k*−1)
- The last bucket is a catchall for statements that have a latency exceeding previous buckets in the histogram.

- `COUNT_BUCKET`

  The number of statements measured with a latency in the interval from `BUCKET_TIMER_LOW` up to but not including `BUCKET_TIMER_HIGH`.

- `COUNT_BUCKET_AND_LOWER`

  The number of statements measured with a latency in the interval from 0 up to but not including `BUCKET_TIMER_HIGH`.

- `BUCKET_QUANTILE`

  The proportion of statements that fall into this or a lower bucket. This proportion corresponds by definition to `COUNT_BUCKET_AND_LOWER / SUM(COUNT_BUCKET)` and is displayed as a convenience column.

The statement histogram summary tables have these indexes:

- `events_statements_histogram_by_digest`:
  - Unique index on (`SCHEMA_NAME`, `DIGEST`, `BUCKET_NUMBER`)
- `events_statements_histogram_global`:
  - Primary key on (`BUCKET_NUMBER`)

`TRUNCATE TABLE` is permitted for statement histogram summary tables. Truncation sets the `COUNT_BUCKET` and `COUNT_BUCKET_AND_LOWER` columns to 0.

In addition, truncating `events_statements_summary_by_digest` implicitly truncates `events_statements_histogram_by_digest`, and truncating `events_statements_summary_global_by_event_name` implicitly truncates `events_statements_histogram_global`.

25.11.15.5 Transaction Summary Tables

The Performance Schema maintains tables for collecting current and recent transaction events, and aggregates that information in summary tables. Section 25.11.7, “Performance Schema Transaction Tables” describes the events on which transaction summaries are based. See that discussion for information about the content of transaction events, the current and recent transaction event tables, and how to control transaction event collection, which is disabled by default.

Example transaction event summary information:

```
mysql> SELECT * FROM events_transactions_summary_global_by_event_name LIMIT 1\G
*************************** 1. row ***************************
          EVENT_NAME: transaction
          COUNT_STAR: 5
      SUM_TIMER_WAIT: 19550092000
      MIN_TIMER_WAIT: 2954148000
      AVG_TIMER_WAIT: 3910018000
```

```
      MAX_TIMER_WAIT: 5486275000
    COUNT_READ_WRITE: 5
SUM_TIMER_READ_WRITE: 19550092000
MIN_TIMER_READ_WRITE: 2954148000
AVG_TIMER_READ_WRITE: 3910018000
MAX_TIMER_READ_WRITE: 5486275000
     COUNT_READ_ONLY: 0
 SUM_TIMER_READ_ONLY: 0
 MIN_TIMER_READ_ONLY: 0
 AVG_TIMER_READ_ONLY: 0
 MAX_TIMER_READ_ONLY: 0
```

Each transaction summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `events_transactions_summary_by_account_by_event_name` has `USER`, `HOST`, and `EVENT_NAME` columns. Each row summarizes events for a given account (user and host combination) and event name.
- `events_transactions_summary_by_host_by_event_name` has `HOST` and `EVENT_NAME` columns. Each row summarizes events for a given host and event name.
- `events_transactions_summary_by_thread_by_event_name` has `THREAD_ID` and `EVENT_NAME` columns. Each row summarizes events for a given thread and event name.
- `events_transactions_summary_by_user_by_event_name` has `USER` and `EVENT_NAME` columns. Each row summarizes events for a given user and event name.
- `events_transactions_summary_global_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.

Each transaction summary table has these summary columns containing aggregated values:

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns are analogous to the columns of the same names in the wait event summary tables (see Section 25.11.15.1, “Wait Event Summary Tables”), except that the transaction summary tables aggregate events from `events_transactions_current` rather than `events_waits_current`. These columns summarize read-write and read-only transactions.

- `COUNT_READ_WRITE`, `SUM_TIMER_READ_WRITE`, `MIN_TIMER_READ_WRITE`, `AVG_TIMER_READ_WRITE`, `MAX_TIMER_READ_WRITE`

  These are similar to the `COUNT_STAR` and *`xxx`*`_TIMER_WAIT` columns, but summarize read-write transactions only.

- `COUNT_READ_ONLY`, `SUM_TIMER_READ_ONLY`, `MIN_TIMER_READ_ONLY`, `AVG_TIMER_READ_ONLY`, `MAX_TIMER_READ_ONLY`

  These are similar to the `COUNT_STAR` and *`xxx`*`_TIMER_WAIT` columns, but summarize read-only transactions only.

The transaction summary tables have these indexes:

- `events_transactions_summary_by_account_by_event_name`:
  - Primary key on (`USER`, `HOST`, `EVENT_NAME`)
- `events_transactions_summary_by_host_by_event_name`:
  - Primary key on (`HOST`, `EVENT_NAME`)
- `events_transactions_summary_by_thread_by_event_name`:

  - Primary key on (THREAD_ID, EVENT_NAME)

- events_transactions_summary_by_user_by_event_name:
  - Primary key on (USER, EVENT_NAME)
- events_transactions_summary_global_by_event_name:
  - Primary key on (EVENT_NAME)

TRUNCATE TABLE is permitted for transaction summary tables. It has these effects:

- For summary tables not aggregated by account, host, or user, truncation resets the summary columns to zero rather than removing rows.
- For summary tables aggregated by account, host, or user, truncation removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero for the remaining rows.

In addition, each transaction summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of events_transactions_summary_global_by_event_name. For details, see Section 25.11.8, "Performance Schema Connection Tables".

Transaction Aggregation Rules

Transaction event collection occurs without regard to isolation level, access mode, or autocommit mode.

Read-write transactions are generally more resource intensive than read-only transactions, therefore transaction summary tables include separate aggregate columns for read-write and read-only transactions.

Resource requirements may also vary with transaction isolation level. However, presuming that only one isolation level would be used per server, aggregation by isolation level is not provided.

25.11.15.6 Object Wait Summary Table

The Performance Schema maintains the objects_summary_global_by_type table for aggregating object wait events.

Example object wait event summary information:

```
mysql> SELECT * FROM objects_summary_global_by_type\G
...
*************************** 3. row ***************************
   OBJECT_TYPE: TABLE
 OBJECT_SCHEMA: test
   OBJECT_NAME: t
    COUNT_STAR: 3
SUM_TIMER_WAIT: 263126976
MIN_TIMER_WAIT: 1522272
AVG_TIMER_WAIT: 87708678
MAX_TIMER_WAIT: 258428280
...
*************************** 10. row ***************************
   OBJECT_TYPE: TABLE
 OBJECT_SCHEMA: mysql
   OBJECT_NAME: user
    COUNT_STAR: 14
SUM_TIMER_WAIT: 365567592
MIN_TIMER_WAIT: 1141704
AVG_TIMER_WAIT: 26111769
MAX_TIMER_WAIT: 334783032
...
```

The `objects_summary_global_by_type` table has these grouping columns to indicate how the table aggregates events: `OBJECT_TYPE`, `OBJECT_SCHEMA`, and `OBJECT_NAME`. Each row summarizes events for the given object.

`objects_summary_global_by_type` has the same summary columns as the `events_waits_summary_by_xxx` tables. See Section 25.11.15.1, “Wait Event Summary Tables”.

The `objects_summary_global_by_type` table has these indexes:

- Primary key on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`)

`TRUNCATE TABLE` is permitted for the object summary table. It resets the summary columns to zero rather than removing rows.

25.11.15.7 File I/O Summary Tables

The Performance Schema maintains file I/O summary tables that aggregate information about I/O operations.

Example file I/O event summary information:

```
mysql> SELECT * FROM file_summary_by_event_name\G
...
*************************** 2. row ***************************
              EVENT_NAME: wait/io/file/sql/binlog
              COUNT_STAR: 31
          SUM_TIMER_WAIT: 8243784888
          MIN_TIMER_WAIT: 0
          AVG_TIMER_WAIT: 265928484
          MAX_TIMER_WAIT: 6490658832
...
mysql> SELECT * FROM file_summary_by_instance\G
...
*************************** 2. row ***************************
               FILE_NAME: /var/mysql/share/english/errmsg.sys
              EVENT_NAME: wait/io/file/sql/ERRMSG
              EVENT_NAME: wait/io/file/sql/ERRMSG
    OBJECT_INSTANCE_BEGIN: 4686193384
              COUNT_STAR: 5
          SUM_TIMER_WAIT: 13990154448
          MIN_TIMER_WAIT: 26349624
          AVG_TIMER_WAIT: 2798030607
          MAX_TIMER_WAIT: 8150662536
...
```

Each file I/O summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `file_summary_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.
- `file_summary_by_instance` has `FILE_NAME`, `EVENT_NAME`, and `OBJECT_INSTANCE_BEGIN` columns. Each row summarizes events for a given file and event name.

Each file I/O summary table has the following summary columns containing aggregated values. Some columns are more general and have values that are the same as the sum of the values of more fine-grained columns. In this way, aggregations at higher levels are available directly without the need for user-defined views that sum lower-level columns.

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns aggregate all I/O operations.

- `COUNT_READ`, `SUM_TIMER_READ`, `MIN_TIMER_READ`, `AVG_TIMER_READ`, `MAX_TIMER_READ`, `SUM_NUMBER_OF_BYTES_READ`

These columns aggregate all read operations, including `FGETS`, `FGETC`, `FREAD`, and `READ`.

- `COUNT_WRITE`, `SUM_TIMER_WRITE`, `MIN_TIMER_WRITE`, `AVG_TIMER_WRITE`, `MAX_TIMER_WRITE`, `SUM_NUMBER_OF_BYTES_WRITE`

  These columns aggregate all write operations, including `FPUTS`, `FPUTC`, `FPRINTF`, `VFPRINTF`, `FWRITE`, and `PWRITE`.

- `COUNT_MISC`, `SUM_TIMER_MISC`, `MIN_TIMER_MISC`, `AVG_TIMER_MISC`, `MAX_TIMER_MISC`

  These columns aggregate all other I/O operations, including `CREATE`, `DELETE`, `OPEN`, `CLOSE`, `STREAM_OPEN`, `STREAM_CLOSE`, `SEEK`, `TELL`, `FLUSH`, `STAT`, `FSTAT`, `CHSIZE`, `RENAME`, and `SYNC`. There are no byte counts for these operations.

The file I/O summary tables have these indexes:

- `file_summary_by_event_name`:
  - Primary key on (`EVENT_NAME`)
- `file_summary_by_instance`:
  - Primary key on (`OBJECT_INSTANCE_BEGIN`)
  - Index on (`FILE_NAME`)
  - Index on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for file I/O summary tables. It resets the summary columns to zero rather than removing rows.

The MySQL server uses several techniques to avoid I/O operations by caching information read from files, so it is possible that statements you might expect to result in I/O events will not. You may be able to ensure that I/O does occur by flushing caches or restarting the server to reset its state.

25.11.15.8 Table I/O and Lock Wait Summary Tables

The following sections describe the table I/O and lock wait summary tables:

- `table_io_waits_summary_by_index_usage`: Table I/O waits per index
- `table_io_waits_summary_by_table`: Table I/O waits per table
- `table_lock_waits_summary_by_table`: Table lock waits per table

The table_io_waits_summary_by_table Table

The `table_io_waits_summary_by_table` table aggregates all table I/O wait events, as generated by the `wait/io/table/sql/handler` instrument. The grouping is by table.

The `table_io_waits_summary_by_table` table has these grouping columns to indicate how the table aggregates events: `OBJECT_TYPE`, `OBJECT_SCHEMA`, and `OBJECT_NAME`. These columns have the same meaning as in the `events_waits_current` table. They identify the table to which the row applies.

`table_io_waits_summary_by_table` has the following summary columns containing aggregated values. As indicated in the column descriptions, some columns are more general and have values that are the same as the sum of the values of more fine-grained columns. For example, columns that aggregate all writes hold the sum of the corresponding columns that aggregate inserts, updates, and deletes. In this way, aggregations at higher levels are available directly without the need for user-defined views that sum lower-level columns.

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns aggregate all I/O operations. They are the same as the sum of the corresponding *`xxx`*`_READ` and *`xxx`*`_WRITE` columns.

- `COUNT_READ`, `SUM_TIMER_READ`, `MIN_TIMER_READ`, `AVG_TIMER_READ`, `MAX_TIMER_READ`

  These columns aggregate all read operations. They are the same as the sum of the corresponding *`xxx`*`_FETCH` columns.

- `COUNT_WRITE`, `SUM_TIMER_WRITE`, `MIN_TIMER_WRITE`, `AVG_TIMER_WRITE`, `MAX_TIMER_WRITE`

  These columns aggregate all write operations. They are the same as the sum of the corresponding *`xxx`*`_INSERT`, *`xxx`*`_UPDATE`, and *`xxx`*`_DELETE` columns.

- `COUNT_FETCH`, `SUM_TIMER_FETCH`, `MIN_TIMER_FETCH`, `AVG_TIMER_FETCH`, `MAX_TIMER_FETCH`

  These columns aggregate all fetch operations.

- `COUNT_INSERT`, `SUM_TIMER_INSERT`, `MIN_TIMER_INSERT`, `AVG_TIMER_INSERT`, `MAX_TIMER_INSERT`

  These columns aggregate all insert operations.

- `COUNT_UPDATE`, `SUM_TIMER_UPDATE`, `MIN_TIMER_UPDATE`, `AVG_TIMER_UPDATE`, `MAX_TIMER_UPDATE`

  These columns aggregate all update operations.

- `COUNT_DELETE`, `SUM_TIMER_DELETE`, `MIN_TIMER_DELETE`, `AVG_TIMER_DELETE`, `MAX_TIMER_DELETE`

  These columns aggregate all delete operations.

The `table_io_waits_summary_by_table` table has these indexes:

- Unique index on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`)

`TRUNCATE TABLE` is permitted for table I/O summary tables. It resets the summary columns to zero rather than removing rows. Truncating this table also truncates the `table_io_waits_summary_by_index_usage` table.

The table_io_waits_summary_by_index_usage Table

The `table_io_waits_summary_by_index_usage` table aggregates all table index I/O wait events, as generated by the `wait/io/table/sql/handler` instrument. The grouping is by table index.

The columns of `table_io_waits_summary_by_index_usage` are nearly identical to `table_io_waits_summary_by_table`. The only difference is the additional group column, `INDEX_NAME`, which corresponds to the name of the index that was used when the table I/O wait event was recorded:

- A value of `PRIMARY` indicates that table I/O used the primary index.
- A value of `NULL` means that table I/O used no index.
- Inserts are counted against `INDEX_NAME = NULL`.

The `table_io_waits_summary_by_index_usage` table has these indexes:

- Unique index on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`, `INDEX_NAME`)

`TRUNCATE TABLE` is permitted for table I/O summary tables. It resets the summary columns to zero rather than removing rows. This table is also truncated by truncation of the `table_io_waits_summary_by_table` table. A DDL operation that changes the index structure of a table may cause the per-index statistics to be reset.

The table_lock_waits_summary_by_table Table

The `table_lock_waits_summary_by_table` table aggregates all table lock wait events, as generated by the `wait/lock/table/sql/handler` instrument. The grouping is by table.

This table contains information about internal and external locks:

- An internal lock corresponds to a lock in the SQL layer. This is currently implemented by a call to `thr_lock()`. In event rows, these locks are distinguished by the `OPERATION` column, which has one of these values:

```
read normal
read with shared locks
read high priority
read no insert
write allow write
write concurrent insert
write delayed
write low priority
write normal
```

- An external lock corresponds to a lock in the storage engine layer. This is currently implemented by a call to `handler::external_lock()`. In event rows, these locks are distinguished by the `OPERATION` column, which has one of these values:

```
read external
write external
```

The `table_lock_waits_summary_by_table` table has these grouping columns to indicate how the table aggregates events: `OBJECT_TYPE`, `OBJECT_SCHEMA`, and `OBJECT_NAME`. These columns have the same meaning as in the `events_waits_current` table. They identify the table to which the row applies.

`table_lock_waits_summary_by_table` has the following summary columns containing aggregated values. As indicated in the column descriptions, some columns are more general and have values that are the same as the sum of the values of more fine-grained columns. For example, columns that aggregate all locks hold the sum of the corresponding columns that aggregate read and write locks. In this way, aggregations at higher levels are available directly without the need for user-defined views that sum lower-level columns.

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns aggregate all lock operations. They are the same as the sum of the corresponding *`xxx`*`_READ` and *`xxx`*`_WRITE` columns.

- `COUNT_READ`, `SUM_TIMER_READ`, `MIN_TIMER_READ`, `AVG_TIMER_READ`, `MAX_TIMER_READ`

  These columns aggregate all read-lock operations. They are the same as the sum of the corresponding *`xxx`*`_READ_NORMAL`, *`xxx`*`_READ_WITH_SHARED_LOCKS`, *`xxx`*`_READ_HIGH_PRIORITY`, and *`xxx`*`_READ_NO_INSERT` columns.

- `COUNT_WRITE`, `SUM_TIMER_WRITE`, `MIN_TIMER_WRITE`, `AVG_TIMER_WRITE`, `MAX_TIMER_WRITE`

  These columns aggregate all write-lock operations. They are the same as the sum of the corresponding *`xxx`*`_WRITE_ALLOW_WRITE`, *`xxx`*`_WRITE_CONCURRENT_INSERT`, *`xxx`*`_WRITE_LOW_PRIORITY`, and *`xxx`*`_WRITE_NORMAL` columns.

- `COUNT_READ_NORMAL`, `SUM_TIMER_READ_NORMAL`, `MIN_TIMER_READ_NORMAL`, `AVG_TIMER_READ_NORMAL`, `MAX_TIMER_READ_NORMAL`

  These columns aggregate internal read locks.

- `COUNT_READ_WITH_SHARED_LOCKS`, `SUM_TIMER_READ_WITH_SHARED_LOCKS`, `MIN_TIMER_READ_WITH_SHARED_LOCKS`, `AVG_TIMER_READ_WITH_SHARED_LOCKS`, `MAX_TIMER_READ_WITH_SHARED_LOCKS`

  These columns aggregate internal read locks.

- `COUNT_READ_HIGH_PRIORITY`, `SUM_TIMER_READ_HIGH_PRIORITY`, `MIN_TIMER_READ_HIGH_PRIORITY`, `AVG_TIMER_READ_HIGH_PRIORITY`, `MAX_TIMER_READ_HIGH_PRIORITY`

  These columns aggregate internal read locks.

- `COUNT_READ_NO_INSERT`, `SUM_TIMER_READ_NO_INSERT`, `MIN_TIMER_READ_NO_INSERT`, `AVG_TIMER_READ_NO_INSERT`, `MAX_TIMER_READ_NO_INSERT`

  These columns aggregate internal read locks.

- `COUNT_READ_EXTERNAL`, `SUM_TIMER_READ_EXTERNAL`, `MIN_TIMER_READ_EXTERNAL`, `AVG_TIMER_READ_EXTERNAL`, `MAX_TIMER_READ_EXTERNAL`

  These columns aggregate external read locks.

- `COUNT_WRITE_ALLOW_WRITE`, `SUM_TIMER_WRITE_ALLOW_WRITE`, `MIN_TIMER_WRITE_ALLOW_WRITE`, `AVG_TIMER_WRITE_ALLOW_WRITE`, `MAX_TIMER_WRITE_ALLOW_WRITE`

  These columns aggregate internal write locks.

- `COUNT_WRITE_CONCURRENT_INSERT`, `SUM_TIMER_WRITE_CONCURRENT_INSERT`, `MIN_TIMER_WRITE_CONCURRENT_INSERT`, `AVG_TIMER_WRITE_CONCURRENT_INSERT`, `MAX_TIMER_WRITE_CONCURRENT_INSERT`

  These columns aggregate internal write locks.

- `COUNT_WRITE_LOW_PRIORITY`, `SUM_TIMER_WRITE_LOW_PRIORITY`, `MIN_TIMER_WRITE_LOW_PRIORITY`, `AVG_TIMER_WRITE_LOW_PRIORITY`, `MAX_TIMER_WRITE_LOW_PRIORITY`

  These columns aggregate internal write locks.

- `COUNT_WRITE_NORMAL`, `SUM_TIMER_WRITE_NORMAL`, `MIN_TIMER_WRITE_NORMAL`, `AVG_TIMER_WRITE_NORMAL`, `MAX_TIMER_WRITE_NORMAL`

  These columns aggregate internal write locks.

- `COUNT_WRITE_EXTERNAL`, `SUM_TIMER_WRITE_EXTERNAL`, `MIN_TIMER_WRITE_EXTERNAL`, `AVG_TIMER_WRITE_EXTERNAL`, `MAX_TIMER_WRITE_EXTERNAL`

  These columns aggregate external write locks.

The `table_lock_waits_summary_by_table` table has these indexes:

- Unique index on (`OBJECT_TYPE`, `OBJECT_SCHEMA`, `OBJECT_NAME`)

`TRUNCATE TABLE` is permitted for table lock summary tables. It resets the summary columns to zero rather than removing rows.

25.11.15.9 Socket Summary Tables

These socket summary tables aggregate timer and byte count information for socket operations:

- `socket_summary_by_event_name`: Aggregate timer and byte count statistics generated by the `wait/io/socket/*` instruments for all socket I/O operations, per socket instrument.
- `socket_summary_by_instance`: Aggregate timer and byte count statistics generated by the `wait/io/socket/*` instruments for all socket I/O operations, per socket instance. When a connection terminates, the row in `socket_summary_by_instance` corresponding to it is deleted.

The socket summary tables do not aggregate waits generated by `idle` events while sockets are waiting for the next request from the client. For `idle` event aggregations, use the wait-event summary tables; see Section 25.11.15.1, “Wait Event Summary Tables”.

Each socket summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `socket_summary_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.
- `socket_summary_by_instance` has an `OBJECT_INSTANCE_BEGIN` column. Each row summarizes events for a given object.

Each socket summary table has these summary columns containing aggregated values:

- `COUNT_STAR`, `SUM_TIMER_WAIT`, `MIN_TIMER_WAIT`, `AVG_TIMER_WAIT`, `MAX_TIMER_WAIT`

  These columns aggregate all operations.

- `COUNT_READ`, `SUM_TIMER_READ`, `MIN_TIMER_READ`, `AVG_TIMER_READ`, `MAX_TIMER_READ`, `SUM_NUMBER_OF_BYTES_READ`

  These columns aggregate all receive operations (`RECV`, `RECVFROM`, and `RECVMSG`).

- `COUNT_WRITE`, `SUM_TIMER_WRITE`, `MIN_TIMER_WRITE`, `AVG_TIMER_WRITE`, `MAX_TIMER_WRITE`, `SUM_NUMBER_OF_BYTES_WRITE`

  These columns aggregate all send operations (`SEND`, `SENDTO`, and `SENDMSG`).

- `COUNT_MISC`, `SUM_TIMER_MISC`, `MIN_TIMER_MISC`, `AVG_TIMER_MISC`, `MAX_TIMER_MISC`

  These columns aggregate all other socket operations, such as `CONNECT`, `LISTEN`, `ACCEPT`, `CLOSE`, and `SHUTDOWN`. There are no byte counts for these operations.

The `socket_summary_by_instance` table also has an `EVENT_NAME` column that indicates the class of the socket: `client_connection`, `server_tcpip_socket`, `server_unix_socket`. This column can be grouped on to isolate, for example, client activity from that of the server listening sockets.

The socket summary tables have these indexes:

- `socket_summary_by_event_name`:
  - Primary key on (`EVENT_NAME`)
- `socket_summary_by_instance`:
  - Primary key on (`OBJECT_INSTANCE_BEGIN`)
  - Index on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for socket summary tables. Except for `events_statements_summary_by_digest`, tt resets the summary columns to zero rather than removing rows.

25.11.15.10 Memory Summary Tables

The Performance Schema instruments memory usage and aggregates memory usage statistics, detailed by these factors:

- Type of memory used (various caches, internal buffers, and so forth)
- Thread, account, user, host indirectly performing the memory operation

The Performance Schema instruments the following aspects of memory use

- Memory sizes used
- Operation counts
- Low and high water marks

Memory sizes help to understand or tune the memory consumption of the server.

Operation counts help to understand or tune the overall pressure the server is putting on the memory allocator, which has an impact on performance. Allocating a single byte one million times is not the same as allocating one million bytes a single time; tracking both sizes and counts can expose the difference.

Low and high water marks are critical to detect workload spikes, overall workload stability, and possible memory leaks.

Memory summary tables do not contain timing information because memory events are not timed.

For information about collecting memory usage data, see Memory Instrumentation Behavior.

Example memory event summary information:

```
mysql> SELECT * FROM memory_summary_global_by_event_name
       WHERE EVENT_NAME = 'memory/sql/TABLE'\G
*************************** 1. row ***************************
                 EVENT_NAME: memory/sql/TABLE
                COUNT_ALLOC: 1381
                 COUNT_FREE: 924
   SUM_NUMBER_OF_BYTES_ALLOC: 2059873
    SUM_NUMBER_OF_BYTES_FREE: 1407432
              LOW_COUNT_USED: 0
          CURRENT_COUNT_USED: 457
             HIGH_COUNT_USED: 461
    LOW_NUMBER_OF_BYTES_USED: 0
CURRENT_NUMBER_OF_BYTES_USED: 652441
   HIGH_NUMBER_OF_BYTES_USED: 669269
```

Each memory summary table has one or more grouping columns to indicate how the table aggregates events. Event names refer to names of event instruments in the `setup_instruments` table:

- `memory_summary_by_account_by_event_name` has `USER`, `HOST`, and `EVENT_NAME` columns. Each row summarizes events for a given account (user and host combination) and event name.
- `memory_summary_by_host_by_event_name` has `HOST` and `EVENT_NAME` columns. Each row summarizes events for a given host and event name.
- `memory_summary_by_thread_by_event_name` has `THREAD_ID` and `EVENT_NAME` columns. Each row summarizes events for a given thread and event name.
- `memory_summary_by_user_by_event_name` has `USER` and `EVENT_NAME` columns. Each row summarizes events for a given user and event name.
- `memory_summary_global_by_event_name` has an `EVENT_NAME` column. Each row summarizes events for a given event name.

Each memory summary table has these summary columns containing aggregated values:

- `COUNT_ALLOC`, `COUNT_FREE`

  The aggregated numbers of calls to memory-allocation and memory-free functions.

- `SUM_NUMBER_OF_BYTES_ALLOC`, `SUM_NUMBER_OF_BYTES_FREE`

  The aggregated sizes of allocated and freed memory blocks.

- `CURRENT_COUNT_USED`

  The aggregated number of currently allocated blocks that have not been freed yet. This is a convenience column, equal to `COUNT_ALLOC` − `COUNT_FREE`.

- `CURRENT_NUMBER_OF_BYTES_USED`

  The aggregated size of currently allocated memory blocks that have not been freed yet. This is a convenience column, equal to `SUM_NUMBER_OF_BYTES_ALLOC` − `SUM_NUMBER_OF_BYTES_FREE`.

- `LOW_COUNT_USED`, `HIGH_COUNT_USED`

  The low and high water marks corresponding to the `CURRENT_COUNT_USED` column.

- `LOW_NUMBER_OF_BYTES_USED`, `HIGH_NUMBER_OF_BYTES_USED`

  The low and high water marks corresponding to the `CURRENT_NUMBER_OF_BYTES_USED` column.

The memory summary tables have these indexes:

- `memory_summary_by_account_by_event_name`:
  - Primary key on (`USER`, `HOST`, `EVENT_NAME`)
- `memory_summary_by_host_by_event_name`:
  - Primary key on (`HOST`, `EVENT_NAME`)
- `memory_summary_by_thread_by_event_name`:
  - Primary key on (`THREAD_ID`, `EVENT_NAME`)
- `memory_summary_by_user_by_event_name`:
  - Primary key on (`USER`, `EVENT_NAME`)
- `memory_summary_global_by_event_name`:
  - Primary key on (`EVENT_NAME`)

`TRUNCATE TABLE` is permitted for memory summary tables. It has these effects:

- In general, truncation resets the baseline for statistics, but does not change the server state. That is, truncating a memory table does not free memory.
- `COUNT_ALLOC` and `COUNT_FREE` are reset to a new baseline, by reducing each counter by the same value.
- Likewise, `SUM_NUMBER_OF_BYTES_ALLOC` and `SUM_NUMBER_OF_BYTES_FREE` are reset to a new baseline.
- `LOW_COUNT_USED` and `HIGH_COUNT_USED` are reset to `CURRENT_COUNT_USED`.
- `LOW_NUMBER_OF_BYTES_USED` and `HIGH_NUMBER_OF_BYTES_USED` are reset to `CURRENT_NUMBER_OF_BYTES_USED`.

In addition, each memory summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of `memory_summary_global_by_event_name`. For details, see Section 25.11.8, “Performance Schema Connection Tables”.

Memory Instrumentation Behavior

Memory instruments are listed in the `setup_instruments` table and have names of the form `memory/code_area/instrument_name`. Memory instrumentation is enabled by default. (Prior to 8.0.2, most memory instrumentation is disabled by default.)

Instruments named with the prefix `memory/performance_schema/` expose how much memory is allocated for internal buffers in the Performance Schema itself. The `memory/performance_schema/` instruments are built in, always enabled, and cannot be disabled at startup or runtime. Built-in memory instruments are displayed only in the `memory_summary_global_by_event_name` table.

To control memory instrumention state at server startup, use lines like these in your `my.cnf` file:

- Enable:

```
[mysqld]
performance-schema-instrument='memory/%=ON'
```

- Disable:

```
[mysqld]
performance-schema-instrument='memory/%=OFF'
```

To control memory instrumention state at runtime, update the `ENABLED` column of the relevant instruments in the `setup_instruments` table:

- Enable:

```
UPDATE setup_instruments SET ENABLED = 'YES'
WHERE NAME LIKE 'memory/%';
```

- Disable:

```
UPDATE setup_instruments SET ENABLED = 'NO'
WHERE NAME LIKE 'memory/%';
```

For memory instruments, the `TIMED` column in `setup_instruments` is ignored because memory operations are not timed.

Note

If disabled memory instrumentation is later enabled, accounting may result in an apparently negative allocation from missing an allocation operation but counting a free operation.

When a thread in the server executes a memory allocation that has been instrumented, these rules apply:

- If the thread is not instrumented or the memory instrument is not enabled, the memory block allocated is not instrumented.
- Otherwise (that is, both the thread and the instrument are enabled), the memory block allocated is instrumented.

For deallocation, these rules apply:

- If a thread is instrumented, and a memory block is not instrumented, the free operation is not instrumented; no statistics are changed.
- If a thread is not instrumented, and a memory block is instrumented, the free operation is instrumented, and statistics are changed.

For the per-thread statistics, the following rules apply.

When an instrumented memory block of size *N* is allocated, the Performance Schema makes these updates to memory summary table columns:

- `COUNT_ALLOC`: Increased by 1
- `CURRENT_COUNT_USED`: Increased by 1
- `HIGH_COUNT_USED`: Increased if `CURRENT_COUNT_USED` is a new maximum
- `SUM_NUMBER_OF_BYTES_ALLOC`: Increased by *N*
- `CURRENT_NUMBER_OF_BYTES_USED`: Increased by *N*
- `HIGH_NUMBER_OF_BYTES_USED`: Increased if `CURRENT_NUMBER_OF_BYTES_USED` is a new maximum

When an instrumented memory block is deallocated, the Performance Schema makes these updates to memory summary table columns:

- `COUNT_FREE`: Increased by 1
- `CURRENT_COUNT_USED`: Decreased by 1
- `LOW_COUNT_USED`: Decreased if `CURRENT_COUNT_USED` is a new minimum
- `SUM_NUMBER_OF_BYTES_FREE`: Increased by *N*
- `CURRENT_NUMBER_OF_BYTES_USED`: Decreased by *N*
- `LOW_NUMBER_OF_BYTES_USED`: Decreased if `CURRENT_NUMBER_OF_BYTES_USED` is a new minimum

For higher-level aggregates (global, by account, by user, by host), the same rules apply as expected for low and high water marks.

- `LOW_COUNT_USED` and `LOW_NUMBER_OF_BYTES_USED` are lower estimates. The value reported by the Performance Schema is guaranteed to be less than or equal to the lowest count or size of memory effectively used at runtime.
- `HIGH_COUNT_USED` and `HIGH_NUMBER_OF_BYTES_USED` are higher estimates. The value reported by the Performance Schema is guaranteed to be greater than or equal to the highest count or size of memory effectively used at runtime.

For lower estimates in summary tables other than `memory_summary_global_by_event_name`, it is possible for values to go negative if memory ownership is transferred between threads.

Here is an example of estimate computation; but note that estimate implementation is subject to change:

Thread 1 uses memory in the range from 1MB to 2MB during execution, as reported by the `LOW_NUMBER_OF_BYTES_USED` and `HIGH_NUMBER_OF_BYTES_USED` columns of the `memory_summary_by_thread_by_event_name` table.

Thread 2 uses memory in the range from 10MB to 12MB during execution, as reported likewise.

When these two threads belong to the same user account, the per-account summary estimates that this account used memory in the range from 11MB to 14MB. That is, the `LOW_NUMBER_OF_BYTES_USED` for the higher level aggregate is the sum of each `LOW_NUMBER_OF_BYTES_USED` (assuming the worst case). Likewise, the `HIGH_NUMBER_OF_BYTES_USED` for the higher level aggregate is the sum of each `HIGH_NUMBER_OF_BYTES_USED` (assuming the worst case).

11MB is a lower estimate that can occur only if both threads hit the low usage mark at the same time.

14MB is a higher estimate that can occur only if both threads hit the high usage mark at the same time.

The real memory usage for this account could have been in the range from 11.5MB to 13.5MB.

For capacity planning, reporting the worst case is actually the desired behavior, as it shows what can potentially happen when sessions are uncorrelated, which is typically the case.

25.11.15.11 Error Summary Tables

The Performance Schema maintains summary tables for aggregating statistical information about server errors (and warnings). For a list of server errors, see Section B.3, “Server Error Codes and Messages”.

Collection of error information is controlled by the `error` instrument, which is enabled by default. Timing information is not collected.

Each error summary table has three columns that identify the error:

- `ERROR_NUMBER` is the numeric error value. The value is unique.
- `ERROR_NAME` is the symbolic error name corresponding to the `ERROR_NUMBER` value. The value is unique.
- `SQLSTATE` is the SQLSTATE value corresponding to the `ERROR_NUMBER` value. The value is not necessarily unique.

For example, if `ERROR_NUMBER` is 1050, `ERROR_NAME` is `ER_TABLE_EXISTS_ERROR` and `SQLSTATE` is `42S01`.

Example error event summary information:

```
mysql> SELECT * FROM events_errors_summary_global_by_error
       WHERE SUM_ERROR_RAISED <> 0\G
*************************** 1. row ***************************
     ERROR_NUMBER: 1064
       ERROR_NAME: ER_PARSE_ERROR
        SQL_STATE: 42000
 SUM_ERROR_RAISED: 1
SUM_ERROR_HANDLED: 0
       FIRST_SEEN: 2016-06-28 07:34:02
        LAST_SEEN: 2016-06-28 07:34:02
*************************** 2. row ***************************
     ERROR_NUMBER: 1146
       ERROR_NAME: ER_NO_SUCH_TABLE
        SQL_STATE: 42S02
 SUM_ERROR_RAISED: 2
SUM_ERROR_HANDLED: 0
       FIRST_SEEN: 2016-06-28 07:34:05
        LAST_SEEN: 2016-06-28 07:36:18
*************************** 3. row ***************************
     ERROR_NUMBER: 1317
       ERROR_NAME: ER_QUERY_INTERRUPTED
        SQL_STATE: 70100
 SUM_ERROR_RAISED: 1
SUM_ERROR_HANDLED: 0
```

```
       FIRST_SEEN: 2016-06-28 11:01:49
        LAST_SEEN: 2016-06-28 11:01:49
```

Each error summary table has one or more grouping columns to indicate how the table aggregates errors:

- `events_errors_summary_by_account_by_error` has `USER`, `HOST`, and `ERROR_NUMBER` columns. Each row summarizes events for a given account (user and host combination) and error.
- `events_errors_summary_by_host_by_error` has `HOST` and `ERROR_NUMBER` columns. Each row summarizes events for a given host and error.
- `events_errors_summary_by_thread_by_error` has `THREAD_ID` and `ERROR_NUMBER` columns. Each row summarizes events for a given thread and error.
- `events_errors_summary_by_user_by_error` has `USER` and `ERROR_NUMBER` columns. Each row summarizes events for a given user and error.
- `events_errors_summary_global_by_error` has an `ERROR_NUMBER` column. Each row summarizes events for a given error.

Each error summary table has these summary columns containing aggregated values:

- `SUM_ERROR_RAISED`

  This column aggregates the number of times the error occurred.

- `SUM_ERROR_HANDLED`

  This column aggregates the number of times the error was handled by an SQL exception handler.

- `FIRST_SEEN`, `LAST_SEEN`

  Timestamp indicating when the error was first seen and most recently seen.

A `NULL` row in each error summary table is used to aggregate statistics for all errors that lie out of range of the instrumented errors. For example, if MySQL Server errors lie in the range from *`M`* to *`N`* and an error is raised with number *`Q`* not in that range, the error is aggregated in the `NULL` row. The `NULL` row is the row with `ERROR_NUMBER=0`, `ERROR_NAME=NULL`, and `SQLSTATE=NULL`.

The error summary tables have these indexes:

- `events_errors_summary_by_account_by_error`:
  - Primary key on (`USER`, `HOST`, `ERROR_NUMBER`)
- `events_errors_summary_by_host_by_error`:
  - Primary key on (`HOST`, `ERROR_NUMBER`)
- `events_errors_summary_by_thread_by_error`:
  - Primary key on (`THREAD_ID`, `ERROR_NUMBER`)
- `events_errors_summary_by_user_by_error`:
  - Primary key on (`USER`, `ERROR_NUMBER`)
- `events_errors_summary_global_by_error`:
  - Primary key on (`ERROR_NUMBER`)

`TRUNCATE TABLE` is permitted for error summary tables. It has these effects:

- For summary tables not aggregated by account, host, or user, truncation resets the summary columns to zero or `NULL` rather than removing rows.
- For summary tables aggregated by account, host, or user, truncation removes rows for accounts, hosts, or users with no connections, and resets the summary columns to zero or `NULL` for the remaining rows.

In addition, each error summary table that is aggregated by account, host, user, or thread is implicitly truncated by truncation of the connection table on which it depends, or truncation of `events_errors_summary_global_by_error`. For details, see Section 25.11.8, “Performance Schema Connection Tables”.

25.11.15.12 Status Variable Summary Tables

The Performance Schema makes status variable information available in the tables described in Section 25.11.14, “Performance Schema Status Variable Tables”. It also makes aggregated status variable information available in summary tables, described here. Each status variable summary table has one or more grouping columns to indicate how the table aggregates status values:

- `status_by_account` has `USER`, `HOST`, and `VARIABLE_NAME` columns to summarize status variables by account.
- `status_by_host` has `HOST` and `VARIABLE_NAME` columns to summarize status variables by the host from which clients connected.
- `status_by_user` has `USER` and `VARIABLE_NAME` columns to summarize status variables by client user name.

Each status variable summary table has this summary column containing aggregated values:

- `VARIABLE_VALUE`

  The aggregated status variable value for active and terminated sessions.

The status variable summary tables have these indexes:

- `status_by_account`:
  - Primary key on (`USER`, `HOST`, `VARIABLE_NAME`)
- `status_by_host`:
  - Primary key on (`HOST`, `VARIABLE_NAME`)
- `status_by_user`:
  - Primary key on (`USER`, `VARIABLE_NAME`)

The meaning of “account” in these tables is similar to its meaning in the MySQL grant tables in the `mysql` system database, in the sense that the term refers to a combination of user and host values. They differ in that, for grant tables, the host part of an account can be a pattern, whereas for Performance Schema tables, the host value is always a specific nonpattern host name.

Account status is collected when sessions terminate. The session status counters are added to the global status counters and the corresponding account status counters. If account statistics are not collected, the session status is added to host and user status, if host and user status are collected.

Account, host, and user statistics are not collected if the `performance_schema_accounts_size`, `performance_schema_hosts_size`, and `performance_schema_users_size` system variables, respectively, are set to 0.

The Performance Schema supports `TRUNCATE TABLE` for status variable summary tables as follows; in all cases, status for active sessions is unaffected:

- `status_by_account`: Aggregates account status from terminated sessions to user and host status, then resets account status.
- `status_by_host`: Resets aggregated host status from terminated sessions.
- `status_by_user`: Resets aggregated user status from terminated sessions.

`FLUSH STATUS` adds the session status from all active sessions to the global status variables, resets the status of all active sessions, and resets account, host, and user status values aggregated from disconnected sessions.

25.11.16 Performance Schema Miscellaneous Tables

The following sections describe tables that do not fall into the table categories discussed in the preceding sections:

- `host_cache`: Information from the internal host cache
- `performance_timers`: Which event timers are available
- `threads`: Information about server threads

25.11.16.1 The host_cache Table

The `host_cache` table provides access to the contents of the host cache, which contains client host name and IP address information and is used to avoid DNS lookups. (See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”.) The `host_cache` table exposes the contents of the host cache so that it can be examined using `SELECT` statements. The Performance Schema must be enabled or this table is empty.

The `host_cache` table has these columns:

- `IP`

  The IP address of the client that connected to the server, expressed as a string.

- `HOST`

  The resolved DNS host name for that client IP, or `NULL` if the name is unknown.

- `HOST_VALIDATED`

  Whether the IP-to-host name-to-IP DNS resolution was performed successfully for the client IP. If `HOST_VALIDATED` is `YES`, the `HOST` column is used as the host name corresponding to the IP so that calls to DNS can be avoided. While `HOST_VALIDATED` is `NO`, DNS resolution is attempted again for each connect, until it eventually completes with either a valid result or a permanent error. This information enables the server to avoid caching bad or missing host names during temporary DNS failures, which would affect clients forever.

- `SUM_CONNECT_ERRORS`

  The number of connection errors that are deemed “blocking” (assessed against the `max_connect_errors` system variable). Only protocol handshake errors are counted, and only for hosts that passed validation (`HOST_VALIDATED = YES`).

- `COUNT_HOST_BLOCKED_ERRORS`

  The number of connections that were blocked because `SUM_CONNECT_ERRORS` exceeded the value of the `max_connect_errors` system variable.

- `COUNT_NAMEINFO_TRANSIENT_ERRORS`

The number of transient errors during IP-to-host name DNS resolution.

- `COUNT_NAMEINFO_PERMANENT_ERRORS`

  The number of permanent errors during IP-to-host name DNS resolution.

- `COUNT_FORMAT_ERRORS`

  The number of host name format errors. MySQL does not perform matching of `Host` column values in the `mysql.user` table against host names for which one or more of the initial components of the name are entirely numeric, such as `1.2.example.com`. The client IP address is used instead. For the rationale why this type of matching does not occur, see Section 6.2.4, “Specifying Account Names”.

- `COUNT_ADDRINFO_TRANSIENT_ERRORS`

  The number of transient errors during host name-to-IP reverse DNS resolution.

- `COUNT_ADDRINFO_PERMANENT_ERRORS`

  The number of permanent errors during host name-to-IP reverse DNS resolution.

- `COUNT_FCRDNS_ERRORS`

  The number of forward-confirmed reverse DNS errors. These errors occur when IP-to-host name-to-IP DNS resolution produces an IP address that does not match the client originating IP address.

- `COUNT_HOST_ACL_ERRORS`

  The number of errors that occur because no user from the client host can possibly log in. In such cases, the server returns `ER_HOST_NOT_PRIVILEGED` and does not even ask for a user name or password.

- `COUNT_NO_AUTH_PLUGIN_ERRORS`

  The number of errors due to requests for an unavailable authentication plugin. A plugin can be unavailable if, for example, it was never loaded or a load attempt failed.

- `COUNT_AUTH_PLUGIN_ERRORS`

  The number of errors reported by authentication plugins.

  An authentication plugin can report different error codes to indicate the root cause of a failure. Depending on the type of error, one of these columns is incremented: `COUNT_AUTHENTICATION_ERRORS`, `COUNT_AUTH_PLUGIN_ERRORS`, `COUNT_HANDSHAKE_ERRORS`. New return codes are an optional extension to the existing plugin API. Unknown or unexpected plugin errors are counted in the `COUNT_AUTH_PLUGIN_ERRORS` column.

- `COUNT_HANDSHAKE_ERRORS`

  The number of errors detected at the wire protocol level.

- `COUNT_PROXY_USER_ERRORS`

  The number of errors detected when a proxy user A is proxied to another user B who does not exist.

- `COUNT_PROXY_USER_ACL_ERRORS`

  The number of errors detected when a proxy user A is proxied to another user B who does exist but for whom A does not have the `PROXY` privilege.

- `COUNT_AUTHENTICATION_ERRORS`

The number of errors caused by failed authentication.

- `COUNT_SSL_ERRORS`

  The number of errors due to SSL problems.

- `COUNT_MAX_USER_CONNECTIONS_ERRORS`

  The number of errors caused by exceeding per-user connection quotas. See Section 6.3.6, “Setting Account Resource Limits”.

- `COUNT_MAX_USER_CONNECTIONS_PER_HOUR_ERRORS`

  The number of errors caused by exceeding per-user connections-per-hour quotas. See Section 6.3.6, “Setting Account Resource Limits”.

- `COUNT_DEFAULT_DATABASE_ERRORS`

  The number of errors related to the default database. For example, the database did not exist or the user had no privileges for accessing it.

- `COUNT_INIT_CONNECT_ERRORS`

  The number of errors caused by execution failures of statements in the `init_connect` system variable value.

- `COUNT_LOCAL_ERRORS`

  The number of errors local to the server implementation and not related to the network, authentication, or authorization. For example, out-of-memory conditions fall into this category.

- `COUNT_UNKNOWN_ERRORS`

  The number of other, unknown errors not accounted for by other columns in this table. This column is reserved for future use, in case new error conditions must be reported, and if preserving the backward compatibility and table structure of the `host_cache` table is required.

- `FIRST_SEEN`

  The timestamp of the first connection attempt seen from the client in the `IP` column.

- `LAST_SEEN`

  The timestamp of the last connection attempt seen from the client in the `IP` column.

- `FIRST_ERROR_SEEN`

  The timestamp of the first error seen from the client in the `IP` column.

- `LAST_ERROR_SEEN`

  The timestamp of the last error seen from the client in the `IP` column.

The `host_cache` table has these indexes:

- Primary key on (`IP`)
- Index on (`HOST`)

`FLUSH HOSTS` and `TRUNCATE TABLE host_cache` have the same effect: They clear the host cache. This also removes rows from the `host_cache` table (because it is the visible representation of the cache) and unblocks any blocked hosts (see Section B.5.2.5, “Host 'host_name' is blocked”.)

`FLUSH HOSTS` requires the `RELOAD` privilege. `TRUNCATE TABLE` requires the `DROP` privilege for the `host_cache` table.

25.11.16.2 The performance_timers Table

The `performance_timers` table shows which event timers are available:

```
mysql> SELECT * FROM performance_timers;
+-------------+-----------------+------------------+----------------+
| TIMER_NAME  | TIMER_FREQUENCY | TIMER_RESOLUTION | TIMER_OVERHEAD |
+-------------+-----------------+------------------+----------------+
| CYCLE       |      2389029850 |                1 |             72 |
| NANOSECOND  |      1000000000 |                1 |            112 |
| MICROSECOND |         1000000 |                1 |            136 |
| MILLISECOND |            1036 |                1 |            168 |
| TICK        |             105 |                1 |           2416 |
+-------------+-----------------+------------------+----------------+
```

The timers in `setup_timers` that you can use are those that do not have `NULL` in the other columns. If the values associated with a given timer name are `NULL`, that timer is not supported on your platform.

The `performance_timers` table has these columns:

- `TIMER_NAME`

  The name by which to refer to the timer when configuring the `setup_timers` table.

- `TIMER_FREQUENCY`

  The number of timer units per second. For a cycle timer, the frequency is generally related to the CPU speed. For example, on a system with a 2.4GHz processor, the `CYCLE` may be close to 2400000000.

- `TIMER_RESOLUTION`

  Indicates the number of timer units by which timer values increase. If a timer has a resolution of 10, its value increases by 10 each time.

- `TIMER_OVERHEAD`

  The minimal number of cycles of overhead to obtain one timing with the given timer. The Performance Schema determines this value by invoking the timer 20 times during initialization and picking the smallest value. The total overhead really is twice this amount because the instrumentation invokes the timer at the start and end of each event. The timer code is called only for timed events, so this overhead does not apply for nontimed events.

The `performance_timers` table has these indexes:

- None

`TRUNCATE TABLE` is not permitted for the `performance_timers` table.

25.11.16.3 The threads Table

The `threads` table contains a row for each server thread. Each row contains information about a thread and indicates whether monitoring and historical event logging are enabled for it:

```
mysql> SELECT * FROM threads\G
*************************** 1. row ***************************
          THREAD_ID: 1
               NAME: thread/sql/main
               TYPE: BACKGROUND
     PROCESSLIST_ID: NULL
   PROCESSLIST_USER: NULL
```

```
   PROCESSLIST_HOST: NULL
     PROCESSLIST_DB: NULL
PROCESSLIST_COMMAND: NULL
   PROCESSLIST_TIME: 80284
  PROCESSLIST_STATE: NULL
   PROCESSLIST_INFO: NULL
   PARENT_THREAD_ID: NULL
               ROLE: NULL
       INSTRUMENTED: YES
            HISTORY: YES
    CONNECTION_TYPE: NULL
       THREAD_OS_ID: 489803
     RESOURCE_GROUP: NULL
...
*************************** 4. row ***************************
          THREAD_ID: 51
               NAME: thread/sql/one_connection
               TYPE: FOREGROUND
     PROCESSLIST_ID: 34
   PROCESSLIST_USER: isabella
   PROCESSLIST_HOST: localhost
     PROCESSLIST_DB: performance_schema
PROCESSLIST_COMMAND: Query
   PROCESSLIST_TIME: 0
  PROCESSLIST_STATE: Sending data
   PROCESSLIST_INFO: SELECT * FROM threads
   PARENT_THREAD_ID: 1
               ROLE: NULL
       INSTRUMENTED: YES
            HISTORY: YES
    CONNECTION_TYPE: SSL/TLS
       THREAD_OS_ID: 755399
     RESOURCE_GROUP: NULL
...
```

When the Performance Schema initializes, it populates the `threads` table based on the threads in existence then. Thereafter, a new row is added each time the server creates a thread.

The `INSTRUMENTED` and `HISTORY` column values for new threads are determined by the contents of the `setup_actors` table. For information about how to use the `setup_actors` table to control these columns, see Section 25.4.6, “Pre-Filtering by Thread”.

Removal of rows from the `threads` table occurs when threads end. For a thread associated with a client session, removal occurs when the session ends. If a client has auto-reconnect enabled and the session reconnects after a disconnect, the session becomes associated with a new row in the `threads` table that has a different `PROCESSLIST_ID` value. The initial `INSTRUMENTED` and `HISTORY` values for the new thread may be different from those of the original thread: The `setup_actors` table may have changed in the meantime, and if the `INSTRUMENTED` or `HISTORY` value for the original thread was changed after the row was initialized, the change does not carry over to the new thread.

The `threads` table columns with names having a prefix of `PROCESSLIST_` provide information similar to that available from the `INFORMATION_SCHEMA.PROCESSLIST` table or the `SHOW PROCESSLIST` statement. Thus, all three sources provide thread-monitoring information. Use of `threads` differs from use of the other two sources in these ways:

- Access to `threads` does not require a mutex and has minimal impact on server performance. `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST` have negative performance consequences because they require a mutex.
- `threads` provides additional information for each thread, such as whether it is a foreground or background thread, and the location within the server associated with the thread.
- `threads` provides information about background threads, so it can be used to monitor activity the other thread information sources cannot.
- You can enable or disable thread monitoring (that is, whether events executed by the thread are instrumented) and historical event logging. To control the initial `INSTRUMENTED` and `HISTORY`

values for new foreground threads, use the `setup_actors` table. To control these aspects of existing threads, set the `INSTRUMENTED` and `HISTORY` columns of `threads` table rows. (For more information about the conditions under which thread monitoring and historical event logging occur, see the descriptions of the `INSTRUMENTED` and `HISTORY` columns.)

For these reasons, DBAs who perform server monitoring using `INFORMATION_SCHEMA.PROCESSLIST` or `SHOW PROCESSLIST` may wish to monitor using the `threads` table instead.

Note

For `INFORMATION_SCHEMA.PROCESSLIST` and `SHOW PROCESSLIST`, information about threads for other users is shown only if the current user has the `PROCESS` privilege. That is not true of the `threads` table; all rows are shown to any user who has the `SELECT` privilege for the table. Users who should not be able to see threads for other users should not be given that privilege.

The `threads` table has these columns:

- `THREAD_ID`

  A unique thread identifier.

- `NAME`

  The name associated with the thread instrumentation code in the server. For example, `thread/sql/one_connection` corresponds to the thread function in the code responsible for handling a user connection, and `thread/sql/main` stands for the `main()` function of the server.

- `TYPE`

  The thread type, either `FOREGROUND` or `BACKGROUND`. User connection threads are foreground threads. Threads associated with internal server activity are background threads. Examples are internal `InnoDB` threads, “binlog dump” threads sending information to slaves, and slave I/O and SQL threads.

- `PROCESSLIST_ID`

  For threads that are displayed in the `INFORMATION_SCHEMA.PROCESSLIST` table, this is the same value displayed in the `ID` column of that table. It is also the value displayed in the `Id` column of `SHOW PROCESSLIST` output, and the value that `CONNECTION_ID()` would return within that thread.

  For background threads (threads not associated with a user connection), `PROCESSLIST_ID` is `NULL`, so the values are not unique.

- `PROCESSLIST_USER`

  The user associated with a foreground thread, `NULL` for a background thread.

- `PROCESSLIST_HOST`

  The host name of the client associated with a foreground thread, `NULL` for a background thread.

  Unlike the `HOST` column of the `INFORMATION_SCHEMA PROCESSLIST` table or the `Host` column of `SHOW PROCESSLIST` output, the `PROCESSLIST_HOST` column does not include the port number for TCP/IP connections. To obtain this information from the Performance Schema, enable the socket instrumentation (which is not enabled by default) and examine the `socket_instances` table:

```
mysql> SELECT * FROM setup_instruments WHERE NAME LIKE 'wait/io/socket%';
+----------------------------------------+---------+-------+
| NAME                                   | ENABLED | TIMED |
```

```
+---------------------------------------+---------+-------+
| wait/io/socket/sql/server_tcpip_socket | NO      | NO    |
| wait/io/socket/sql/server_unix_socket  | NO      | NO    |
| wait/io/socket/sql/client_connection   | NO      | NO    |
+---------------------------------------+---------+-------+
3 rows in set (0.01 sec)

mysql> UPDATE setup_instruments SET ENABLED='YES' WHERE NAME LIKE 'wait/io/socket%';
Query OK, 3 rows affected (0.00 sec)
Rows matched: 3  Changed: 3  Warnings: 0

mysql> SELECT * FROM socket_instances\G
*************************** 1. row ***************************
          EVENT_NAME: wait/io/socket/sql/client_connection
OBJECT_INSTANCE_BEGIN: 140612577298432
            THREAD_ID: 31
            SOCKET_ID: 53
                   IP: ::ffff:127.0.0.1
                 PORT: 55642
                STATE: ACTIVE
...
```

- `PROCESSLIST_DB`

  The default database for the thread, or `NULL` if there is none.

- `PROCESSLIST_COMMAND`

  For foreground threads, the type of command the thread is executing on behalf of the client, or `Sleep` if the session is idle. For descriptions of thread commands, see Section 8.14, “Examining Thread Information”. The value of this column corresponds to the `COM_xxx` commands of the client/server protocol and `Com_xxx` status variables. See Section 5.1.7, “Server Status Variables”

  Background threads do not execute commands on behalf of clients, so this column may be `NULL`.

- `PROCESSLIST_TIME`

  The time in seconds that the thread has been in its current state.

- `PROCESSLIST_STATE`

  An action, event, or state that indicates what the thread is doing. For descriptions of `PROCESSLIST_STATE` values, see Section 8.14, “Examining Thread Information”. If the value if `NULL`, the thread may correspond to an idle client session or the work it is doing is not instrumented with stages.

  Most states correspond to very quick operations. If a thread stays in a given state for many seconds, there might be a problem that bears investigation.

- `PROCESSLIST_INFO`

  The statement the thread is executing, or `NULL` if it is not executing any statement. The statement might be the one sent to the server, or an innermost statement if the statement executes other statements. For example, if a `CALL` statement executes a stored procedure that is executing a `SELECT` statement, the `PROCESSLIST_INFO` value shows the `SELECT` statement.

- `PARENT_THREAD_ID`

  If this thread is a subthread (spawned by another thread), this is the `THREAD_ID` value of the spawning thread.

- `ROLE`

  Unused.

- `INSTRUMENTED`

Whether events executed by the thread are instrumented. The value is `YES` or `NO`.

- For foreground threads, the initial `INSTRUMENTED` value is determined by whether the user account associated with the thread matches any row in the `setup_actors` table. Matching is based on the values of the `PROCESSLIST_USER` and `PROCESSLIST_HOST` columns.

  If the thread spawns a subthread, matching occurs again for the `threads` table row created for the subthread.

- For background threads, `INSTRUMENTED` is `YES` by default. `setup_actors` is not consulted because there is no associated user for background threads.

- For any thread, its `INSTRUMENTED` value can be changed during the lifetime of the thread.

For monitoring of events executed by the thread to occur, these things must be true:

- The `thread_instrumentation` consumer in the `setup_consumers` table must be `YES`.

- The `threads.INSTRUMENTED` column must be `YES`.

- Monitoring occurs only for those thread events produced from instruments that have the `ENABLED` column set to `YES` in the `setup_instruments` table.

- `HISTORY`

  Whether to log historical events for the thread. The value is `YES` or `NO`.

  - For foreground threads, the initial `HISTORY` value is determined by whether the user account associated with the thread matches any row in the `setup_actors` table. Matching is based on the values of the `PROCESSLIST_USER` and `PROCESSLIST_HOST` columns.

    If the thread spawns a subthread, matching occurs again for the `threads` table row created for the subthread.

  - For background threads, `HISTORY` is `YES` by default. `setup_actors` is not consulted because there is no associated user for background threads.

  - For any thread, its `HISTORY` value can be changed during the lifetime of the thread.

  For historical event logging for the thread to occur, these things must be true:

  - The appropriate history-related consumers in the `setup_consumers` table must be enabled. For example, wait event logging in the `events_waits_history` and `events_waits_history_long` tables requires the corresponding `events_waits_history` and `events_waits_history_long` consumers to be `YES`.

  - The `threads.HISTORY` column must be `YES`.

  - Logging occurs only for those thread events produced from instruments that have the `ENABLED` column set to `YES` in the `setup_instruments` table.

- `CONNECTION_TYPE`

  The protocol used to establish the connection, or `NULL` for background threads. Permitted values are `TCP/IP` (TCP/IP connection established without SSL), `SSL/TLS` (TCP/IP connection established with SSL), `Socket` (Unix socket file connection), `Named Pipe` (Windows named pipe connection), and `Shared Memory` (Windows shared memory connection).

- `THREAD_OS_ID`

  The thread or task identifier as defined by the underlying operating system, if there is one:

- When a MySQL thread is associated with the same operating system thread for its lifetime, `THREAD_OS_ID` contains the operating system thread ID.
- When a MySQL thread is not associated with the same operating system thread for its lifetime, `THREAD_OS_ID` contains `NULL`. This is typical for user sessions when the thread pool plugin is used (see MySQL Enterprise Thread Pool).

For Windows, `THREAD_OS_ID` corresponds to the thread ID visible in Process Explorer (https://technet.microsoft.com/en-us/sysinternals/bb896653.aspx).

For Linux, `THREAD_OS_ID` corresponds to the value of the `gettid()` function. This value is exposed, for example, using the `perf` or `ps -L` commands, or in the `proc` file system (`/proc/[pid]/task/[tid]`). For more information, see the `perf-stat(1)`, `ps(1)`, and `proc(5)` man pages.

- `RESOURCE_GROUP`

  The resource group label, or `NULL` if there is no label.

  This column was added in MySQL 8.0.2

The `threads` table has these indexes:

- Primary key on (`THREAD_ID`)
- Index on (`NAME`)
- Index on (`PROCESSLIST_ID`)
- Index on (`PROCESSLIST_USER`, `PROCESSLIST_HOST`)
- Index on (`PROCESSLIST_HOST`)
- Index on (`THREAD_OS_ID`)
- Index on (`RESOURCE_GROUP`); added in MySQL 8.0.2

`TRUNCATE TABLE` is not permitted for the `threads` table.

25.11.16.4 The user_defined_functions Table

The `user_defined_functions` table contains a row for each user-defined function (UDF) registered by a server component or plugin, or by a `CREATE FUNCTION` statement. This table was added in MySQL 8.0.2.

`user_defined_functions` has these columns:

- `UDF_NAME`

  The UDF name. The value is `NULL` if the function was registered by a `CREATE FUNCTION` statement and is in the process of unloading.

- `UDF_RETURN_TYPE`

  The UDF return value type. The value is one of `int`, `decimal`, `real`, `char`, or `row`.

- `UDF_TYPE`

  The UDF type. The value is one of `function` (scalar) or `aggregate`.

- `UDF_LIBRARY`

  The UDF library file. The value is `NULL` if the UDF was registered using a server component service rather than by a `CREATE FUNCTION` statement.

- `UDF_USAGE_COUNT`

  The current UDF usage count.

The `user_defined_functions` table has these indexes:

- Primary key on (`UDF_NAME`)

`TRUNCATE TABLE` is not permitted for the `user_defined_functions` table.

25.12 Performance Schema Option and Variable Reference

Table 25.3 Performance Schema Variable Reference

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
performance_schema	Yes	Yes	Yes		Global	No
Performance_schema_accounts_lost				Yes	Global	No
performance_schema_accounts_size	Yes	Yes	Yes		Global	No
Performance_schema_cond_classes_lost				Yes	Global	No
Performance_schema_cond_instances_lost				Yes	Global	No
performance-schema-consumer-events-stages-current	Yes	Yes				
performance-schema-consumer-events-stages-history	Yes	Yes				
performance-schema-consumer-events-stages-history-long	Yes	Yes				
performance-schema-consumer-events-statements-current	Yes	Yes				
performance-schema-consumer-events-statements-history	Yes	Yes				
performance-schema-consumer-events-statements-history-long	Yes	Yes				
performance-schema-	Yes	Yes				

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
consumer-events-transactions-current						
performance-schema-consumer-events-transactions-history	Yes	Yes				
performance-schema-consumer-events-transactions-history-long	Yes	Yes				
performance-schema-consumer-events-waits-current	Yes	Yes				
performance-schema-consumer-events-waits-history	Yes	Yes				
performance-schema-consumer-events-waits-history-long	Yes	Yes				
performance-schema-consumer-global-instrumentation	Yes	Yes				
performance-schema-consumer-statements-digest	Yes	Yes				
performance-schema-consumer-thread-instrumentation	Yes	Yes				
Performance_schema_digest_lost				Yes	Global	No
performance_schema_digests_size	Yes	Yes	Yes		Global	No
performance_schema_events_stages_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_stages_history_size	Yes	Yes	Yes		Global	No
performance_schema_events_statements_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_statements_history_size	Yes	Yes	Yes		Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
performance_schema_events_transactions_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_transactions_history_size	Yes	Yes	Yes		Global	No
performance_schema_events_waits_history_long_size	Yes	Yes	Yes		Global	No
performance_schema_events_waits_history_size	Yes	Yes	Yes		Global	No
Performance_schema_file_classes_lost				Yes	Global	No
Performance_schema_file_handles_lost				Yes	Global	No
Performance_schema_file_instances_lost				Yes	Global	No
Performance_schema_hosts_lost				Yes	Global	No
performance_schema_hosts_size	Yes	Yes	Yes		Global	No
performance-schema-instrument	Yes	Yes				
Performance_schema_locker_lost				Yes	Global	No
performance_schema_max_cond_classes	Yes	Yes	Yes		Global	No
performance_schema_max_cond_instances	Yes	Yes	Yes		Global	No
performance_schema_max_digest_length	Yes	Yes	Yes		Global	No
performance_schema_max_file_classes	Yes	Yes	Yes		Global	No
performance_schema_max_file_handles	Yes	Yes	Yes		Global	No
performance_schema_max_file_instances	Yes	Yes	Yes		Global	No
performance_schema_max_memory_classes	Yes	Yes	Yes		Global	No
performance_schema_max_metadata_locks	Yes	Yes	Yes		Global	No
performance_schema_max_mutex_classes	Yes	Yes	Yes		Global	No
performance_schema_max_mutex_instances	Yes	Yes	Yes		Global	No
performance_schema_max_prepared_statements_instances	Yes	Yes	Yes		Global	No
performance_schema_max_program_instances	Yes	Yes	Yes		Global	No
performance_schema_max_rwlock_classes	Yes	Yes	Yes		Global	No
performance_schema_max_rwlock_instances	Yes	Yes	Yes		Global	No
performance_schema_max_socket_classes	Yes	Yes	Yes		Global	No
performance_schema_max_socket_instances	Yes	Yes	Yes		Global	No
performance_schema_max_stage_classes	Yes	Yes	Yes		Global	No
performance_schema_max_statement_classes	Yes	Yes	Yes		Global	No
performance_schema_max_statement_stack	Yes	Yes	Yes		Global	No
performance_schema_max_table_handles	Yes	Yes	Yes		Global	No
performance_schema_max_table_instances	Yes	Yes	Yes		Global	No
performance_schema_max_thread_classes	Yes	Yes	Yes		Global	No
performance_schema_max_thread_instances	Yes	Yes	Yes		Global	No
Performance_schema_memory_classes_lost				Yes	Global	No
Performance_schema_metadata_lock_lost				Yes	Global	No
Performance_schema_mutex_classes_lost				Yes	Global	No
Performance_schema_mutex_instances_lost				Yes	Global	No
Performance_schema_nested_statement_lost				Yes	Global	No

Name	Cmd-Line	Option File	System Var	Status Var	Var Scope	Dynamic
Performance_schema_prepared_statements_lost				Yes	Global	No
Performance_schema_program_lost				Yes	Global	No
Performance_schema_rwlock_classes_lost				Yes	Global	No
Performance_schema_rwlock_instances_lost				Yes	Global	No
Performance_schema_session_connect_attrs_lost				Yes	Global	No
performance_schema_session_connect_attrs_size	Yes	Yes	Yes		Global	No
performance_schema_setup_actors_size	Yes	Yes	Yes		Global	No
performance_schema_setup_objects_size	Yes	Yes	Yes		Global	No
Performance_schema_socket_classes_lost				Yes	Global	No
Performance_schema_socket_instances_lost				Yes	Global	No
Performance_schema_stage_classes_lost				Yes	Global	No
Performance_schema_statement_classes_lost				Yes	Global	No
Performance_schema_table_handles_lost				Yes	Global	No
Performance_schema_table_instances_lost				Yes	Global	No
Performance_schema_thread_classes_lost				Yes	Global	No
Performance_schema_thread_instances_lost				Yes	Global	No
Performance_schema_users_lost				Yes	Global	No
performance_schema_users_size	Yes	Yes	Yes		Global	No

25.13 Performance Schema Command Options

Performance Schema parameters can be specified at server startup on the command line or in option files to configure Performance Schema instruments and consumers. Runtime configuration is also possible in many cases (see Section 25.4, “Performance Schema Runtime Configuration”), but startup configuration must be used when runtime configuration is too late to affect instruments that have already been initialized during the startup process.

Performance Schema consumers and instruments can be configured at startup using the following syntax. For additional details, see Section 25.3, “Performance Schema Startup Configuration”.

- `--performance-schema-consumer-`*`consumer_name`*`=value`

  Configure a Performance Schema consumer. Consumer names in the `setup_consumers` table use underscores, but for consumers set at startup, dashes and underscores within the name are equivalent. Options for configuring individual consumers are detailed later in this section.

- `--performance-schema-instrument=`*`instrument_name`*`=value`

  Configure a Performance Schema instrument. The name may be given as a pattern to configure instruments that match the pattern.

The following items configure individual consumers:

- `--performance-schema-consumer-events-stages-current=value`

  Configure the `events-stages-current` consumer.

- `--performance-schema-consumer-events-stages-history=value`

  Configure the `events-stages-history` consumer.

- `--performance-schema-consumer-events-stages-history-long=value`

Configure the `events-stages-history-long` consumer.

- `--performance-schema-consumer-events-statements-current=value`

  Configure the `events-statements-current` consumer.

- `--performance-schema-consumer-events-statements-history=value`

  Configure the `events-statements-history` consumer.

- `--performance-schema-consumer-events-statements-history-long=value`

  Configure the `events-statements-history-long` consumer.

- `--performance-schema-consumer-events-transactions-current=value`

  Configure the Performance Schema `events-transactions-current` consumer.

- `--performance-schema-consumer-events-transactions-history=value`

  Configure the Performance Schema `events-transactions-history` consumer.

- `--performance-schema-consumer-events-transactions-history-long=value`

  Configure the Performance Schema `events-transactions-history-long` consumer.

- `--performance-schema-consumer-events-waits-current=value`

  Configure the `events-waits-current` consumer.

- `--performance-schema-consumer-events-waits-history=value`

  Configure the `events-waits-history` consumer.

- `--performance-schema-consumer-events-waits-history-long=value`

  Configure the `events-waits-history-long` consumer.

- `--performance-schema-consumer-global-instrumentation=value`

  Configure the `global-instrumentation` consumer.

- `--performance-schema-consumer-statements-digest=value`

  Configure the `statements-digest` consumer.

- `--performance-schema-consumer-thread-instrumentation=value`

  Configure the `thread-instrumentation` consumer.

25.14 Performance Schema System Variables

The Performance Schema implements several system variables that provide configuration information:

```
mysql> SHOW VARIABLES LIKE 'perf%';
+-----------------------------------------------------+-------+
| Variable_name                                       | Value |
+-----------------------------------------------------+-------+
| performance_schema                                  | ON    |
| performance_schema_accounts_size                    | -1    |
| performance_schema_digests_size                     | 10000 |
| performance_schema_events_stages_history_long_size  | 10000 |
| performance_schema_events_stages_history_size       | 10    |
| performance_schema_events_statements_history_long_size | 10000 |
```

```
| performance_schema_events_statements_history_size         | 10    |
| performance_schema_events_transactions_history_long_size | 10000 |
| performance_schema_events_transactions_history_size       | 10    |
| performance_schema_events_waits_history_long_size         | 10000 |
| performance_schema_events_waits_history_size              | 10    |
| performance_schema_hosts_size                             | -1    |
| performance_schema_max_cond_classes                       | 80    |
| performance_schema_max_cond_instances                     | -1    |
| performance_schema_max_digest_length                      | 1024  |
| performance_schema_max_file_classes                       | 50    |
| performance_schema_max_file_handles                       | 32768 |
| performance_schema_max_file_instances                     | -1    |
| performance_schema_max_index_stat                         | -1    |
| performance_schema_max_memory_classes                     | 320   |
| performance_schema_max_metadata_locks                     | -1    |
| performance_schema_max_mutex_classes                      | 220   |
| performance_schema_max_mutex_instances                    | -1    |
| performance_schema_max_prepared_statements_instances      | -1    |
| performance_schema_max_program_instances                  | -1    |
| performance_schema_max_rwlock_classes                     | 40    |
| performance_schema_max_rwlock_instances                   | -1    |
| performance_schema_max_socket_classes                     | 10    |
| performance_schema_max_socket_instances                   | -1    |
| performance_schema_max_sql_text_length                    | 1024  |
| performance_schema_max_stage_classes                      | 150   |
| performance_schema_max_statement_classes                  | 192   |
| performance_schema_max_statement_stack                    | 10    |
| performance_schema_max_table_handles                      | -1    |
| performance_schema_max_table_instances                    | -1    |
| performance_schema_max_table_lock_stat                    | -1    |
| performance_schema_max_thread_classes                     | 50    |
| performance_schema_max_thread_instances                   | -1    |
| performance_schema_session_connect_attrs_size             | 512   |
| performance_schema_setup_actors_size                      | -1    |
| performance_schema_setup_objects_size                     | -1    |
| performance_schema_users_size                             | -1    |
+-----------------------------------------------------------+-------+
```

Performance Schema system variables can be set at server startup on the command line or in option files, and many can be set at runtime. See Section 25.12, “Performance Schema Option and Variable Reference”.

The Performance Schema automatically sizes the values of several of its parameters at server startup if they are not set explicitly. For more information, see Section 25.3, “Performance Schema Startup Configuration”.

Performance Schema system variables have the following meanings:

- `performance_schema`

Command-Line Format	`--performance-schema=#`	
System Variable	**Name**	`performance_schema`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	boolean
	Default	`ON`

The value of this variable is `ON` or `OFF` to indicate whether the Performance Schema is enabled. By default, the value is `ON`. At server startup, you can specify this variable with no value or a value of `ON` or 1 to enable it, or with a value of `OFF` or 0 to disable it.

Even when the Performance Schema is disabled, it continues to populate the `global_variables`, `session_variables`, `global_status`, and `session_status` tables. This occurs as necessary

to permit the results for the `SHOW VARIABLES` and `SHOW STATUS` statements to be drawn from those tables, depending on the setting of the `show_compatibiliy_56` system variable.

- `performance_schema_accounts_size`

Command-Line Format	`--performance-schema-accounts-size=#`	
System Variable	Name	`performance_schema_accounts_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autoscaled)`
	Min Value	`-1 (autoscaled)`
	Max Value	`1048576`

The number of rows in the `accounts` table. If this variable is 0, the Performance Schema does not maintain connection statistics in the `accounts` table or status variable information in the `status_by_account` table.

- `performance_schema_digests_size`

Command-Line Format	`--performance-schema-digests-size=#`	
System Variable	Name	`performance_schema_digests_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autosized)`
	Min Value	`-1`
	Max Value	`1048576`

The maximum number of rows in the `events_statements_summary_by_digest` table. If this maximum is exceeded such that a digest cannot be instrumented, the Performance Schema increments the `Performance_schema_digest_lost` status variable.

For more information about statement digesting, see Section 25.9, “Performance Schema Statement Digests”.

- `performance_schema_error_size`

Command-Line Format	`--performance-schema-error-size`	
System Variable	Name	`performance_schema_error_size`
	Variable Scope	Global
	Dynamic Variable	No

Permitted Values	**Type**	integer
	Default	`number of server error codes`
	Min Value	`0`
	Max Value	`1048576`

The number of instrumented server error codes. The default value is the actual number of server error codes. Although the value can be set anywhere from 0 to its maximum, the intended use is to set it to either its default (to instrument all errors) or 0 (to instrument no errors).

Error information is aggregated in summary tables; see Section 25.11.15.11, “Error Summary Tables”. If an error occurs that is not instrumented, information for the occurrence is aggregated to the `NULL` row in each summary table; that is, to the row with `ERROR_NUMBER=0`, `ERROR_NAME=NULL`, and `SQLSTATE=NULL`.

- `performance_schema_events_stages_history_long_size`

Command-Line Format	`--performance-schema-events-stages-history-long-size=#`	
System Variable	**Name**	`performance_schema_events_stages_history_long_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows in the `events_stages_history_long` table.

- `performance_schema_events_stages_history_size`

Command-Line Format	`--performance-schema-events-stages-history-size=#`	
System Variable	**Name**	`performance_schema_events_stages_history_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows per thread in the `events_stages_history` table.

- `performance_schema_events_statements_history_long_size`

Command-Line Format	`--performance-schema-events-statements-history-long-size=#`	
System Variable	**Name**	`performance_schema_events_statements_history_long_siz`
	Variable Scope	Global
	Dynamic Variable	No

Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows in the `events_statements_history_long` table.

- `performance_schema_events_statements_history_size`

Command-Line Format	`--performance-schema-events-statements-history-size=#`	
System Variable	**Name**	`performance_schema_events_statements_history_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows per thread in the `events_statements_history` table.

- `performance_schema_events_transactions_history_long_size`

Command-Line Format	`--performance-schema-events-transactions-history-long-size=#`	
System Variable	**Name**	`performance_schema_events_transactions_history_long_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows in the `events_transactions_history_long` table.

- `performance_schema_events_transactions_history_size`

Command-Line Format	`--performance-schema-events-transactions-history-size=#`	
System Variable	**Name**	`performance_schema_events_transactions_history_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows per thread in the `events_transactions_history` table.

- `performance_schema_events_waits_history_long_size`

Command-Line Format	`--performance-schema-events-waits-history-long-size=#`	
System Variable	**Name**	`performance_schema_events_waits_history_long_size`
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows in the `events_waits_history_long` table.

- `performance_schema_events_waits_history_size`

Command-Line Format	`--performance-schema-events-waits-history-size=#`	
System Variable	**Name**	`performance_schema_events_waits_history_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The number of rows per thread in the `events_waits_history` table.

- `performance_schema_hosts_size`

Command-Line Format	`--performance-schema-hosts-size=#`	
System Variable	**Name**	`performance_schema_hosts_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`
	Min Value	`-1 (autoscaled)`
	Max Value	`1048576`

The number of rows in the `hosts` table. If this variable is 0, the Performance Schema does not maintain connection statistics in the `hosts` table or status variable information in the `status_by_host` table.

- `performance_schema_max_cond_classes`

Command-Line Format	`--performance-schema-max-cond-classes=#`	
System Variable	**Name**	`performance_schema_max_cond_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`80`

The maximum number of condition instruments.

- `performance_schema_max_cond_instances`

Command-Line Format	`--performance-schema-max-cond-instances=#`	
System Variable	Name	`performance_schema_max_cond_instances`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented condition objects.

- `performance_schema_max_digest_length`

Command-Line Format	`--performance-schema-max-digest-length=#`	
System Variable	Name	`performance_schema_max_digest_length`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`1024`
	Min Value	`0`
	Max Value	`1048576`

The maximum number of bytes available for storage of normalized statement digest values in the Performance Schema. This variable is related to `max_digest_length`; see the description of that variable in Section 5.1.5, “Server System Variables”

For more information about statement digesting, see Section 25.9, “Performance Schema Statement Digests”.

- `performance_schema_max_file_classes`

Command-Line Format	`--performance-schema-max-file-classes=#`	
System Variable	Name	`performance_schema_max_file_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`80`

The maximum number of file instruments.

- `performance_schema_max_file_handles`

Command-Line Format	`--performance-schema-max-file-handles=#`	
System Variable	Name	`performance_schema_max_file_handles`

	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`32768`

The maximum number of opened file objects.

The value of `performance_schema_max_file_handles` should be greater than the value of `open_files_limit`: `open_files_limit` affects the maximum number of open file handles the server can support and `performance_schema_max_file_handles` affects how many of these file handles can be instrumented.

- `performance_schema_max_file_instances`

Command-Line Format	`--performance-schema-max-file-instances=#`	
System Variable	**Name**	`performance_schema_max_file_instances`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented file objects.

- `performance_schema_max_index_stat`

Command-Line Format	`--performance-schema-max-index-stat=#`	
System Variable	**Name**	`performance_schema_max_index_stat`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autosized)`

The maximum number of indexes for which the Performance Schema maintains statistics. If this maximum is exceeded such that index statistics are lost, the Performance Schema increments the `Performance_schema_index_stat_lost` status variable. The default value is autosized using the value of `performance_schema_max_table_instances`.

- `performance_schema_max_memory_classes`

Command-Line Format	`--performance-schema-max-memory-classes=#`	
System Variable	**Name**	`performance_schema_max_memory_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer

	Default	`320`

The maximum number of memory instruments.

- `performance_schema_max_metadata_locks`

Command-Line Format	`--performance-schema-max-metadata-locks=#`	
System Variable	**Name**	`performance_schema_max_metadata_locks`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The maximum number of metadata lock instruments. This value controls the size of the `metadata_locks` table. If this maximum is exceeded such that a metadata lock cannot be instrumented, the Performance Schema increments the `Performance_schema_metadata_lock_lost` status variable.

- `performance_schema_max_mutex_classes`

Command-Line Format	`--performance-schema-max-mutex-classes=#`	
System Variable	**Name**	`performance_schema_max_mutex_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values (8.0.0)	**Type**	integer
	Default	`200`
Permitted Values (>= 8.0.1, <= 8.0.2)	**Type**	integer
	Default	`220`
Permitted Values (>= 8.0.3)	**Type**	integer
	Default	`250`

The maximum number of mutex instruments.

- `performance_schema_max_mutex_instances`

Command-Line Format	`--performance-schema-max-mutex-instances=#`	
System Variable	**Name**	`performance_schema_max_mutex_instances`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented mutex objects.

- `performance_schema_max_prepared_statements_instances`

Command-Line Format	--performance-schema-max-prepared-statements-instances=#	
System Variable	Name	performance_schema_max_prepared_statements_instances
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	-1 (autoscaled)

The maximum number of rows in the prepared_statements_instances table. If this maximum is exceeded such that a prepared statement cannot be instrumented, the Performance Schema increments the Performance_schema_prepared_statements_lost status variable. The default value of this variable is autosized based on the value of the max_prepared_stmt_count system variable.

- performance_schema_max_rwlock_classes

Command-Line Format	--performance-schema-max-rwlock-classes=#	
System Variable	Name	performance_schema_max_rwlock_classes
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	40

The maximum number of rwlock instruments.

- performance_schema_max_program_instances

Command-Line Format	--performance-schema-max-program-instances=#	
System Variable	Name	performance_schema_max_program_instances
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	-1 (autoscaled)

The maximum number of stored programs for which the Performance Schema maintains statistics. If this maximum is exceeded, the Performance Schema increments the Performance_schema_program_lost status variable.

- performance_schema_max_rwlock_instances

Command-Line Format	--performance-schema-max-rwlock-instances=#	
System Variable	Name	performance_schema_max_rwlock_instances
	Variable Scope	Global

	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented rwlock objects.

- `performance_schema_max_socket_classes`

Command-Line Format	`--performance-schema-max-socket-classes=#`	
System Variable	**Name**	`performance_schema_max_socket_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`10`

The maximum number of socket instruments.

- `performance_schema_max_socket_instances`

Command-Line Format	`--performance-schema-max-socket-instances=#`	
System Variable	**Name**	`performance_schema_max_socket_instances`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented socket objects.

- `performance_schema_max_sql_text_length`

Command-Line Format	`--performance-schema-max-sql-text-length=#`	
System Variable	**Name**	`performance_schema_max_sql_text_length`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`1024`
	Min Value	`0`
	Max Value	`1048576`

The maximum number of bytes used to store SQL statements in the `SQL_TEXT` column of the `events_statements_current`, `events_statements_history`, and `events_statements_history_long` statement event tables. Any bytes in excess of

`performance_schema_max_sql_text_length` are discarded and do not appear in the `SQL_TEXT` column. Statements differing only after that many initial bytes are indistinguishable in this column.

Decreasing the `performance_schema_max_sql_text_length` value reduces memory use but causes more statements to become indistinguishable if they differ only at the end. Increasing the value increases memory use but permits longer statements to be distinguished.

- `performance_schema_max_stage_classes`

Command-Line Format	`--performance-schema-max-stage-classes=#`	
System Variable	Name	`performance_schema_max_stage_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`150`

The maximum number of stage instruments.

- `performance_schema_max_statement_classes`

Command-Line Format	`--performance-schema-max-statement-classes=#`	
System Variable	Name	`performance_schema_max_statement_classes`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autosized)`

The maximum number of statement instruments. The default value is calculated at server build time based on the number of commands in the client/server protocol and the number of SQL statement types supported by the server.

This variable should not be changed, unless to set it to 0 to disable all statement instrumentation and save all memory associated with it. Setting the variable to nonzero values other than the default has no benefit; in particular, values larger than the default cause more memory to be allocated then is needed.

- `performance_schema_max_statement_stack`

Command-Line Format	`--performance-schema-max-statement-stack=#`	
System Variable	Name	`performance_schema_max_statement_stack`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`10`

The maximum depth of nested stored program calls for which the Performance Schema maintains statistics. When this maximum is exceeded, the Performance Schema increments the `Performance_schema_nested_statement_lost` status variable for each stored program statement executed.

- `performance_schema_max_table_handles`

Command-Line Format	`--performance-schema-max-table-handles=#`	
System Variable	Name	`performance_schema_max_table_handles`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autoscaled)`

The maximum number of opened table objects. This value controls the size of the `table_handles` table. If this maximum is exceeded such that a table handle cannot be instrumented, the Performance Schema increments the `Performance_schema_table_handles_lost` status variable.

- `performance_schema_max_table_instances`

Command-Line Format	`--performance-schema-max-table-instances=#`	
System Variable	Name	`performance_schema_max_table_instances`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autoscaled)`

The maximum number of instrumented table objects.

- `performance_schema_max_table_lock_stat`

Command-Line Format	`--performance-schema-max-table-lock-stat=#`	
System Variable	Name	`performance_schema_max_table_lock_stat`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	`-1 (autosized)`

The maximum number of tables for which the Performance Schema maintains lock statistics. If this maximum is exceeded such that table lock statistics are lost, the Performance Schema increments the `Performance_schema_table_lock_stat_lost` status variable.

- `performance_schema_max_thread_classes`

Command-Line Format	`--performance-schema-max-thread-classes=#`

System Variable	Name	performance_schema_max_thread_classes
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	50

The maximum number of thread instruments.

- performance_schema_max_thread_instances

Command-Line Format	--performance-schema-max-thread-instances=#	
System Variable	Name	performance_schema_max_thread_instances
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	-1 (autoscaled)

The maximum number of instrumented thread objects. The value controls the size of the threads table. If this maximum is exceeded such that a thread cannot be instrumented, the Performance Schema increments the Performance_schema_thread_instances_lost status variable.

The max_connections system variable affects how many threads can run in the server. performance_schema_max_thread_instances affects how many of these running threads can be instrumented.

The variables_by_thread and status_by_thread tables contain system and status variable information information only about foreground threads. If not all threads are instrumented by the Performance Schema, this table will miss some rows. In this case, the Performance_schema_thread_instances_lost status variable will be greater than zero.

- performance_schema_session_connect_attrs_size

Command-Line Format	--performance-schema-session-connect-attrs-size=#	
System Variable	Name	performance_schema_session_connect_attrs_size
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	Type	integer
	Default	-1 (autosized)
	Min Value	-1
	Max Value	1048576

The amount of preallocated memory per thread reserved to hold connection attribute key/value pairs. If the aggregate size of connection attribute data sent by a client is larger than this amount, the Performance Schema truncates the attribute data, increments the

`Performance_schema_session_connect_attrs_lost` status variable, and writes a message to the error log indicating that truncation occurred if the `log_warnings` system variable is greater than zero. A `_truncated` attribute is also added to the session attributes with a value indicating how many bytes were lost, if the attribute buffer has sufficient space. This enables the Performance Schema to expose per-connection truncation information in the connection attribute tables. This information can be examined without having to check the error log.

The default value of `performance_schema_session_connect_attrs_size` is autosized at server startup. This value may be small, so if truncation occurs (`Performance_schema_session_connect_attrs_lost` becomes nonzero), you may wish to set `performance_schema_session_connect_attrs_size` explicitly to a larger value.

Although the maximum permitted `performance_schema_session_connect_attrs_size` value is 1MB, the effective maximum is 64KB because the server imposes a limit of 64KB on the aggregate size of connection attribute data it will accept. If a client attempts to send more than 64KB of attribute data, the server rejects the connection. For more information, see Section 25.11.9, “Performance Schema Connection Attribute Tables”.

- `performance_schema_setup_actors_size`

Command-Line Format	`--performance-schema-setup-actors-size=#`	
System Variable	**Name**	`performance_schema_setup_actors_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The number of rows in the `setup_actors` table.

- `performance_schema_setup_objects_size`

Command-Line Format	`--performance-schema-setup-objects-size=#`	
System Variable	**Name**	`performance_schema_setup_objects_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer
	Default	`-1 (autoscaled)`

The number of rows in the `setup_objects` table.

- `performance_schema_users_size`

Command-Line Format	`--performance-schema-users-size=#`	
System Variable	**Name**	`performance_schema_users_size`
	Variable Scope	Global
	Dynamic Variable	No
Permitted Values	**Type**	integer

	Default	`-1 (autoscaled)`
	Min Value	`-1 (autoscaled)`
	Max Value	`1048576`

The number of rows in the `users` table. If this variable is 0, the Performance Schema does not maintain connection statistics in the `users` table or status variable information in the `status_by_user` table.

25.15 Performance Schema Status Variables

The Performance Schema implements several status variables that provide information about instrumentation that could not be loaded or created due to memory constraints:

```
mysql> SHOW STATUS LIKE 'perf%';
+-------------------------------------------+-------+
| Variable_name                             | Value |
+-------------------------------------------+-------+
| Performance_schema_accounts_lost          | 0     |
| Performance_schema_cond_classes_lost      | 0     |
| Performance_schema_cond_instances_lost    | 0     |
| Performance_schema_file_classes_lost      | 0     |
| Performance_schema_file_handles_lost      | 0     |
| Performance_schema_file_instances_lost    | 0     |
| Performance_schema_hosts_lost             | 0     |
| Performance_schema_locker_lost            | 0     |
| Performance_schema_mutex_classes_lost     | 0     |
| Performance_schema_mutex_instances_lost   | 0     |
| Performance_schema_rwlock_classes_lost    | 0     |
| Performance_schema_rwlock_instances_lost  | 0     |
| Performance_schema_socket_classes_lost    | 0     |
| Performance_schema_socket_instances_lost  | 0     |
| Performance_schema_stage_classes_lost     | 0     |
| Performance_schema_statement_classes_lost | 0     |
| Performance_schema_table_handles_lost     | 0     |
| Performance_schema_table_instances_lost   | 0     |
| Performance_schema_thread_classes_lost    | 0     |
| Performance_schema_thread_instances_lost  | 0     |
| Performance_schema_users_lost             | 0     |
+-------------------------------------------+-------+
```

Performance Schema status variables have the following meanings:

- `Performance_schema_accounts_lost`

  The number of times a row could not be added to the `accounts` table because it was full.

- `Performance_schema_cond_classes_lost`

  How many condition instruments could not be loaded.

- `Performance_schema_cond_instances_lost`

  How many condition instrument instances could not be created.

- `Performance_schema_digest_lost`

  The number of digest instances that could not be instrumented in the `events_statements_summary_by_digest` table. This can be nonzero if the value of `performance_schema_digests_size` is too small.

- `Performance_schema_file_classes_lost`

How many file instruments could not be loaded.

- `Performance_schema_file_handles_lost`

  How many file instrument instances could not be opened.

- `Performance_schema_file_instances_lost`

  How many file instrument instances could not be created.

- `Performance_schema_hosts_lost`

  The number of times a row could not be added to the `hosts` table because it was full.

- `Performance_schema_index_stat_lost`

  The number of indexes for which statistics were lost. This can be nonzero if the value of `performance_schema_max_index_stat` is too small.

- `Performance_schema_locker_lost`

  How many events are “lost” or not recorded, due to the following conditions:

  - Events are recursive (for example, waiting for A caused a wait on B, which caused a wait on C).
  - The depth of the nested events stack is greater than the limit imposed by the implementation.

  Events recorded by the Performance Schema are not recursive, so this variable should always be 0.

- `Performance_schema_memory_classes_lost`

  The number of times a memory instrument could not be loaded.

- `Performance_schema_metadata_lock_lost`

  The number of metadata locks that could not be instrumented in the `metadata_locks` table. This can be nonzero if the value of `performance_schema_max_metadata_locks` is too small.

- `Performance_schema_mutex_classes_lost`

  How many mutex instruments could not be loaded.

- `Performance_schema_mutex_instances_lost`

  How many mutex instrument instances could not be created.

- `Performance_schema_nested_statement_lost`

  The number of stored program statements for which statistics were lost. This can be nonzero if the value of `performance_schema_max_statement_stack` is too small.

- `Performance_schema_prepared_statements_lost`

  The number of prepared statements that could not be instrumented in the `prepared_statements_instances` table. This can be nonzero if the value of `performance_schema_max_prepared_statements_instances` is too small.

- `Performance_schema_program_lost`

  The number of stored programs for which statistics were lost. This can be nonzero if the value of `performance_schema_max_program_instances` is too small.

- `Performance_schema_rwlock_classes_lost`

How many rwlock instruments could not be loaded.

- `Performance_schema_rwlock_instances_lost`

  How many rwlock instrument instances could not be created.

- `Performance_schema_session_connect_attrs_longest_seen`

  In addition to the connection attribute size-limit check performed by the Performance Schema against the value of the `performance_schema_session_connect_attrs_size` system variable, the server performs a preliminary check, imposing a limit of 64KB on the aggregate size of connection attribute data it will accept. If a client attempts to send more than 64KB of attribute data, the server rejects the connection. Otherwise, the server considers the attribute buffer valid and tracks the size of the longest such buffer in the `Performance_schema_session_connect_attrs_longest_seen` status variable. If this value is larger than `performance_schema_session_connect_attrs_size`, DBAs may wish to increase the latter value, or, alternatively, investigate which clients are sending large amounts of attribute data.

  For more information about connection attributes, see Section 25.11.9, “Performance Schema Connection Attribute Tables”.

- `Performance_schema_session_connect_attrs_lost`

  The number of connections for which connection attribute truncation has occurred. For a given connection, if the client sends connection attribute key/value pairs for which the aggregate size is larger is larger than the reserved storage permitted by the value of the `performance_schema_session_connect_attrs_size` system variable, the Performance Schema truncates the attribute data and increments `Performance_schema_session_connect_attrs_lost`. If this value is nonzero, you may wish to set `performance_schema_session_connect_attrs_size` to a larger value.

  For more information about connection attributes, see Section 25.11.9, “Performance Schema Connection Attribute Tables”.

- `Performance_schema_socket_classes_lost`

  How many socket instruments could not be loaded.

- `Performance_schema_socket_instances_lost`

  How many socket instrument instances could not be created.

- `Performance_schema_stage_classes_lost`

  How many stage instruments could not be loaded.

- `Performance_schema_statement_classes_lost`

  How many statement instruments could not be loaded.

- `Performance_schema_table_handles_lost`

  How many table instrument instances could not be opened. This can be nonzero if the value of `performance_schema_max_table_handles` is too small.

- `Performance_schema_table_instances_lost`

  How many table instrument instances could not be created.

- `Performance_schema_table_lock_stat_lost`

The number of tables for which lock statistics were lost. This can be nonzero if the value of `performance_schema_max_table_lock_stat` is too small.

- `Performance_schema_thread_classes_lost`

  How many thread instruments could not be loaded.

- `Performance_schema_thread_instances_lost`

  The number of thread instances that could not be instrumented in the `threads` table. This can be nonzero if the value of `performance_schema_max_thread_instances` is too small.

- `Performance_schema_users_lost`

  The number of times a row could not be added to the `users` table because it was full.

For information on using these variables to check Performance Schema status, see Section 25.7, “Performance Schema Status Monitoring”.

25.16 The Performance Schema Memory-Allocation Model

The Performance Schema uses this memory allocation model:

- May allocate memory at server startup
- May allocate additional memory during server operation
- Never free memory during server operation (although it might be recycled)
- Free all memory used at shutdown

The result is to relax memory constraints so that the Performance Schema can be used with less configuration, and to decrease the memory footprint so that consumption scales with server load. Memory used depends on the load actually seen, not the load estimated or explicitly configured for.

Several Performance Schema sizing parameters are autoscaled and need not be configured explicitly unless you want to establish an explicit limit on memory allocation:

```
performance_schema_accounts_size
performance_schema_hosts_size
performance_schema_max_cond_instances
performance_schema_max_file_instances
performance_schema_max_index_stat
performance_schema_max_metadata_locks
performance_schema_max_mutex_instances
performance_schema_max_prepared_statements_instances
performance_schema_max_program_instances
performance_schema_max_rwlock_instances
performance_schema_max_socket_instances
performance_schema_max_table_handles
performance_schema_max_table_instances
performance_schema_max_table_lock_stat
performance_schema_max_thread_instances
performance_schema_users_size
```

For an autoscaled parameter, configuration works like this:

- With the value set to -1 (the default), the parameter is autoscaled:
  - The corresponding internal buffer is empty initially and no memory is allocated.
  - As the Performance Schema collects data, memory is allocated in the corresponding buffer. The buffer size is unbounded, and may grow with the load.
- With the value set to 0:

  - The corresponding internal buffer is empty initially and no memory is allocated.
- With the value set to *N* > 0:
  - The corresponding internal buffer is empty initially and no memory is allocated.
  - As the Performance Schema collects data, memory is allocated in the corresponding buffer, until the buffer size reaches *N*.
  - Once the buffer size reaches *N*, no more memory is allocated. Data collected by the Performance Schema for this buffer is lost, and any corresponding “lost instance” counters are incremented.

To see how much memory the Performance Schema is using, check the instruments designed for that purpose. The Performance Schema allocates memory internally and associates each buffer with a dedicated instrument so that memory consumption can be traced to individual buffers. Instruments named with the prefix `memory/performance_schema/` expose how much memory is allocated for these internal buffers. The buffers are global to the server, so the instruments are displayed only in the `memory_summary_global_by_event_name` table, and not in other `memory_summary_by_xxx_by_event_name` tables.

This query shows the information associated with the memory instruments:

```
SELECT * FROM memory_summary_global_by_event_name
WHERE EVENT_NAME LIKE 'memory/performance_schema/%';
```

25.17 Performance Schema and Plugins

Removing a plugin with `UNINSTALL PLUGIN` does not affect information already collected for code in that plugin. Time spent executing the code while the plugin was loaded was still spent even if the plugin is unloaded later. The associated event information, including aggregate information, remains readable in `performance_schema` database tables. For additional information about the effect of plugin installation and removal, see Section 25.7, “Performance Schema Status Monitoring”.

A plugin implementor who instruments plugin code should document its instrumentation characteristics to enable those who load the plugin to account for its requirements. For example, a third-party storage engine should include in its documentation how much memory the engine needs for mutex and other instruments.

25.18 Using the Performance Schema to Diagnose Problems

The Performance Schema is a tool to help a DBA do performance tuning by taking real measurements instead of “wild guesses.” This section demonstrates some ways to use the Performance Schema for this purpose. The discussion here relies on the use of event filtering, which is described in Section 25.4.2, “Performance Schema Event Filtering”.

The following example provides one methodology that you can use to analyze a repeatable problem, such as investigating a performance bottleneck. To begin, you should have a repeatable use case where performance is deemed “too slow” and needs optimization, and you should enable all instrumentation (no pre-filtering at all).

1. Run the use case.
2. Using the Performance Schema tables, analyze the root cause of the performance problem. This analysis will rely heavily on post-filtering.
3. For problem areas that are ruled out, disable the corresponding instruments. For example, if analysis shows that the issue is not related to file I/O in a particular storage engine, disable the file I/O instruments for that engine. Then truncate the history and summary tables to remove previously collected events.

4. Repeat the process at step 1.

   At each iteration, the Performance Schema output, particularly the `events_waits_history_long` table, will contain less and less “noise” caused by nonsignificant instruments, and given that this table has a fixed size, will contain more and more data relevant to the analysis of the problem at hand.

   At each iteration, investigation should lead closer and closer to the root cause of the problem, as the “signal/noise” ratio will improve, making analysis easier.

5. Once a root cause of performance bottleneck is identified, take the appropriate corrective action, such as:

   - Tune the server parameters (cache sizes, memory, and so forth).
   - Tune a query by writing it differently,
   - Tune the database schema (tables, indexes, and so forth).
   - Tune the code (this applies to storage engine or server developers only).

6. Start again at step 1, to see the effects of the changes on performance.

The `mutex_instances.LOCKED_BY_THREAD_ID` and `rwlock_instances.WRITE_LOCKED_BY_THREAD_ID` columns are extremely important for investigating performance bottlenecks or deadlocks. This is made possible by Performance Schema instrumentation as follows:

1. Suppose that thread 1 is stuck waiting for a mutex.
2. You can determine what the thread is waiting for:

   ```
   SELECT * FROM events_waits_current WHERE THREAD_ID = thread_1;
   ```

   Say the query result identifies that the thread is waiting for mutex A, found in `events_waits_current.OBJECT_INSTANCE_BEGIN`.

3. You can determine which thread is holding mutex A:

   ```
   SELECT * FROM mutex_instances WHERE OBJECT_INSTANCE_BEGIN = mutex_A;
   ```

   Say the query result identifies that it is thread 2 holding mutex A, as found in `mutex_instances.LOCKED_BY_THREAD_ID`.

4. You can see what thread 2 is doing:

   ```
   SELECT * FROM events_waits_current WHERE THREAD_ID = thread_2;
   ```

25.18.1 Query Profiling Using Performance Schema

The following example demonstrates how to use Performance Schema statement events and stage events to retrieve data comparable to profiling information provided by `SHOW PROFILES` and `SHOW PROFILE` statements.

The `setup_actors` table can be used to limit the collection of historical events by host, user, or account to reduce runtime overhead and the amount of data collected in history tables. The first step of the example shows how to limit collection of historical events to a specific user.

Performance Schema displays event timer information in picoseconds (trillionths of a second) to normalize timing data to a standard unit. In the following example, `TIMER_WAIT` values are divided

by 1000000000000 to show data in units of seconds. Values are also truncated to 6 decimal places to display data in the same format as `SHOW PROFILES` and `SHOW PROFILE` statements.

1. Limit the collection of historical events to the user that will run the query. By default, `setup_actors` is configured to allow monitoring and historical event collection for all foreground threads:

```
mysql> SELECT * FROM performance_schema.setup_actors;
+------+------+------+---------+---------+
| HOST | USER | ROLE | ENABLED | HISTORY |
+------+------+------+---------+---------+
| %    | %    | %    | YES     | YES     |
+------+------+------+---------+---------+
```

   Update the default row in the `setup_actors` table to disable historical event collection and monitoring for all foreground threads, and insert a new row that enables monitoring and historical event collection for the user that will run the query:

```
mysql> UPDATE performance_schema.setup_actors SET ENABLED = 'NO', HISTORY = 'NO'
       WHERE HOST = '%' AND USER = '%';

mysql> INSERT INTO performance_schema.setup_actors (HOST,USER,ROLE,ENABLED,HISTORY)
       VALUES('localhost','test_user','%','YES','YES');
```

   Data in the `setup_actors` table should now appear similar to the following:

```
mysql> SELECT * FROM performance_schema.setup_actors;
+-----------+-----------+------+---------+---------+
| HOST      | USER      | ROLE | ENABLED | HISTORY |
+-----------+-----------+------+---------+---------+
| %         | %         | %    | NO      | NO      |
| localhost | test_user | %    | YES     | YES     |
+-----------+-----------+------+---------+---------+
```

2. Ensure that statement and stage instrumentation is enabled by updating the `setup_instruments` table. Some instruments may already be enabled by default.

```
mysql> UPDATE performance_schema.setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
       WHERE NAME LIKE '%statement/%';

mysql> UPDATE performance_schema.setup_instruments SET ENABLED = 'YES', TIMED = 'YES'
       WHERE NAME LIKE '%stage/%';
```

3. Ensure that `events_statements_*` and `events_stages_*` consumers are enabled. Some consumers may already be enabled by default.

```
mysql> UPDATE performance_schema.setup_consumers SET ENABLED = 'YES'
       WHERE NAME LIKE '%events_statements_%';

mysql> UPDATE performance_schema.setup_consumers SET ENABLED = 'YES'
       WHERE NAME LIKE '%events_stages_%';
```

4. Under the user account you are monitoring, run the statement that you want to profile. For example:

```
mysql> SELECT * FROM employees.employees WHERE emp_no = 10001;
+--------+------------+------------+-----------+--------+------------+
| emp_no | birth_date | first_name | last_name | gender | hire_date |
+--------+------------+------------+-----------+--------+------------+
|  10001 | 1953-09-02 | Georgi     | Facello   | M      | 1986-06-26 |
+--------+------------+------------+-----------+--------+------------+
```

5. Identify the `EVENT_ID` of the statement by querying the `events_statements_history_long` table. This step is similar to running `SHOW PROFILES` to identify the `Query_ID`. The following query produces output similar to `SHOW PROFILES`:

```
mysql> SELECT EVENT_ID, TRUNCATE(TIMER_WAIT/1000000000000,6) as Duration, SQL_TEXT
       FROM performance_schema.events_statements_history_long WHERE SQL_TEXT like '%10001%';
+----------+----------+-----------------------------------------------------------+
| event_id | duration | sql_text                                                  |
+----------+----------+-----------------------------------------------------------+
|       31 | 0.028310 | SELECT * FROM employees.employees WHERE emp_no = 10001 |
+----------+----------+-----------------------------------------------------------+
```

6. Query the `events_stages_history_long` table to retrieve the statement's stage events. Stages are linked to statements using event nesting. Each stage event record has a `NESTING_EVENT_ID` column that contains the `EVENT_ID` of the parent statement.

```
mysql> SELECT event_name AS Stage, TRUNCATE(TIMER_WAIT/1000000000000,6) AS Duration
       FROM performance_schema.events_stages_history_long WHERE NESTING_EVENT_ID=31;
+--------------------------------+----------+
| Stage                          | Duration |
+--------------------------------+----------+
| stage/sql/starting             | 0.000080 |
| stage/sql/checking permissions | 0.000005 |
| stage/sql/Opening tables       | 0.027759 |
| stage/sql/init                 | 0.000052 |
| stage/sql/System lock          | 0.000009 |
| stage/sql/optimizing           | 0.000006 |
| stage/sql/statistics           | 0.000082 |
| stage/sql/preparing            | 0.000008 |
| stage/sql/executing            | 0.000000 |
| stage/sql/Sending data         | 0.000017 |
| stage/sql/end                  | 0.000001 |
| stage/sql/query end            | 0.000004 |
| stage/sql/closing tables       | 0.000006 |
| stage/sql/freeing items        | 0.000272 |
| stage/sql/cleaning up          | 0.000001 |
+--------------------------------+----------+
```

Chapter 26 MySQL sys Schema

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MySQL 8.0 includes the `sys` schema, a set of objects that helps DBAs and developers interpret data collected by the Performance Schema. `sys` schema objects can be used for typical tuning and diagnosis use cases. Objects in this schema include:

- Views that summarize Performance Schema data into more easily understandable form.
- Stored procedures that perform operations such as Performance Schema configuration and generating diagnostic reports.
- Stored functions that query Performance Schema configuration and provide formatting services.

For new installations, the `sys` schema is installed by default during data directory initialization if you use `mysqld` with the `--initialize` or `--initialize-insecure` option. If this is not desired, you can drop the `sys` schema manually after initialization if it is unneeded.

`mysql_upgrade` returns an error if a `sys` schema exists but has no `version` view, on the assumption that absence of this view indicates a user-created `sys` schema. To upgrade in this case, remove or rename the existing `sys` schema first.

`sys` schema objects have a `DEFINER` of `'mysql.sys'@'localhost'`. Use of the dedicated `mysql.sys` account avoids problems that occur if a DBA renames or removes the `root` account.

26.1 Prerequisites for Using the sys Schema

Before using the `sys` schema, the prerequisites described in this section must be satisfied.

The `sys` schema requires MySQL 5.6 or higher.

Because the `sys` schema provides an alternative means of accessing the Performance Schema, the Performance Schema must be enabled for the `sys` schema to work. See Section 25.3, “Performance Schema Startup Configuration”.

For full access to the `sys` schema, a user must have these privileges:

- `SELECT` on all `sys` tables and views
- `EXECUTE` on all `sys` stored procedures and functions
- `INSERT` and `UPDATE` for the `sys_config` table, if changes are to be made to it
- Additional privileges for certain `sys` schema stored procedures and functions, as noted in their descriptions; for example, the `ps_setup_save()` procedure

It is also necessary to have privileges for the objects underlying the `sys` schema objects:

- `SELECT` on any Performance Schema tables accessed by `sys` schema objects, and `UPDATE` for any tables to be updated using `sys` schema objects

- `PROCESS` for the `INFORMATION_SCHEMA INNODB_BUFFER_PAGE` table

Certain Performance Schema instruments and consumers must be enabled and (for instruments) timed to take full advantage of `sys` schema capabilities:

- All `wait` instruments
- All `stage` instruments
- All `statement` instruments
- *`xxx`*`_current` and *`xxx`*`_history_long` consumers for all events

You can use the `sys` schema itself to enable all of the additional instruments and consumers:

```
CALL sys.ps_setup_enable_instrument('wait');
CALL sys.ps_setup_enable_instrument('stage');
CALL sys.ps_setup_enable_instrument('statement');
CALL sys.ps_setup_enable_consumer('current');
CALL sys.ps_setup_enable_consumer('history_long');
```

Note

For many uses of the `sys` schema, the default Performance Schema is sufficient for data collection. Enabling all the instruments and consumers just mentioned has a performance impact, so it is preferable to enable only the additional configuration you need. Also, remember that if you enable additional configuration, you can easily restore the default configuration like this:

```
CALL sys.ps_setup_reset_to_default(TRUE);
```

26.2 Using the sys Schema

You can make the `sys` schema the default schema so that references to its objects need not be qualified with the schema name:

```
mysql> USE sys;
Database changed
mysql> SELECT * FROM version;
+-------------+---------------------+
| sys_version | mysql_version       |
+-------------+---------------------+
| 1.6.0       | 8.0.2-dmr-debug-log |
+-------------+---------------------+
```

(The `version` view shows the `sys` schema and MySQL server versions.)

To access `sys` schema objects while a different schema is the default (or simply to be explicit), qualify object references with the schema name:

```
mysql> SELECT * FROM sys.version;
+-------------+---------------------+
| sys_version | mysql_version       |
+-------------+---------------------+
| 1.6.0       | 8.0.2-dmr-debug-log |
+-------------+---------------------+
```

Examples in this chapter usually assume `sys` as the default schema.

The `sys` schema contains many views that summarize Performance Schema tables in various ways. Most of these views come in pairs, such that one member of the pair has the same name as the other

member, plus a `x$` prefix. For example, the `host_summary_by_file_io` view summarizes file I/O grouped by host and displays latencies converted from picoseconds to more readable values (with units);

```
mysql> SELECT * FROM host_summary_by_file_io;
+------------+-------+------------+
| host       | ios   | io_latency |
+------------+-------+------------+
| localhost  | 67570 | 5.38 s     |
| background |  3468 | 4.18 s     |
+------------+-------+------------+
```

The `x$host_summary_by_file_io` view summarizes the same data but displays unformatted picosecond latencies:

```
mysql> SELECT * FROM x$host_summary_by_file_io;
+------------+-------+---------------+
| host       | ios   | io_latency    |
+------------+-------+---------------+
| localhost  | 67574 | 5380678125144 |
| background |  3474 | 4758696829416 |
+------------+-------+---------------+
```

The view without the `x$` prefix is intended to provide output that is more user friendly and easier for humans to read. The view with the `x$` prefix that displays the same values in raw form is intended more for use with other tools that perform their own processing on the data. For additional information about the differences between non-`x$` and `x$` views, see Section 26.4.3, “sys Schema Views”.

To examine `sys` schema object definitions, use the appropriate `SHOW` statement or `INFORMATION_SCHEMA` query. For example, to examine the definitions of the `session` view and `format_bytes()` function, use these statements:

```
mysql> SHOW CREATE VIEW session;
mysql> SHOW CREATE FUNCTION format_bytes;
```

However, those statements display the definitions in relatively unformatted form. To view object definitions with more readable formatting, access the individual `.sql` files available from the `sys` schema development web site at https://github.com/mysql/mysql-sys.

Neither `mysqldump` nor `mysqlpump` dump the `sys` schema by default. To generate a dump file, name the `sys` schema explicitly on the command line using either of these commands:

```
mysqldump --databases --routines sys > sys_dump.sql
mysqlpump sys > sys_dump.sql
```

To reinstall the schema from the dump file, use this command:

```
mysql < sys_dump.sql
```

26.3 sys Schema Progress Reporting

The following `sys` schema views provide progress reporting for long-running transactions:

```
processlist
session
x$processlist
x$session
```

Assuming that the required instruments and consumers are enabled, the `progress` column of these views shows the percentage of work completed for stages that support progress reporting.

Stage progress reporting requires that the `events_stages_current` consumer be enabled, as well as the instruments for which progress information is desired. Instruments for these stages currently support progress reporting:

```
stage/sql/Copying to tmp table
stage/innodb/alter table (end)
stage/innodb/alter table (flush)
stage/innodb/alter table (insert)
stage/innodb/alter table (log apply index)
stage/innodb/alter table (log apply table)
stage/innodb/alter table (merge sort)
stage/innodb/alter table (read PK and internal sort)
stage/innodb/buffer pool load
```

For stages that do not support estimated and completed work reporting, or if the required instruments or consumers are not enabled, the `progress` column is `NULL`.

26.4 sys Schema Object Reference

The `sys` schema includes tables and triggers, views, and stored procedures and functions. The following sections provide details for each of these objects.

26.4.1 sys Schema Object Index

The following tables list `sys` schema objects and provide a short description of each one.

Table 26.1 sys Schema Tables and Triggers

Table or Trigger Name	Description
`sys_config`	`sys` schema configuration options
`sys_config_insert_set_user`	`sys_config` insert trigger
`sys_config_update_set_user`	`sys_config` update trigger

Table 26.2 sys Schema Views

View Name	Description
`host_summary`, `x$host_summary`	Statement activity, file I/O, and connections, grouped by host
`host_summary_by_file_io`, `x$host_summary_by_file_io`	File I/O, grouped by host
`host_summary_by_file_io_type`, `x$host_summary_by_file_io_type`	File I/O, grouped by host and event type
`host_summary_by_stages`, `x$host_summary_by_stages`	Statement stages, grouped by host
`host_summary_by_statement_latency`, `x$host_summary_by_statement_latency`	Statement statistics, grouped by host
`host_summary_by_statement_type`, `x$host_summary_by_statement_type`	Statements executed, grouped by host and statement
`innodb_buffer_stats_by_schema`, `x$innodb_buffer_stats_by_schema`	`InnoDB` buffer information, grouped by schema
`innodb_buffer_stats_by_table`, `x$innodb_buffer_stats_by_table`	`InnoDB` buffer information, grouped by schema and table
`innodb_lock_waits`, `x$innodb_lock_waits`	`InnoDB` lock information
`io_by_thread_by_latency`, `x$io_by_thread_by_latency`	I/O consumers, grouped by thread

View Name	Description
io_global_by_file_by_bytes, x$io_global_by_file_by_bytes	Global I/O consumers, grouped by file and bytes
io_global_by_file_by_latency, x$io_global_by_file_by_latency	Global I/O consumers, grouped by file and latency
io_global_by_wait_by_bytes, x$io_global_by_wait_by_bytes	Global I/O consumers, grouped by bytes
io_global_by_wait_by_latency, x$io_global_by_wait_by_latency	Global I/O consumers, grouped by latency
latest_file_io, x$latest_file_io	Most recent I/O, grouped by file and thread
memory_by_host_by_current_bytes, x$memory_by_host_by_current_bytes	Memory use, grouped by host
memory_by_thread_by_current_bytes, x$memory_by_thread_by_current_bytes	Memory use, grouped by thread
memory_by_user_by_current_bytes, x$memory_by_user_by_current_bytes	Memory use, grouped by user
memory_global_by_current_bytes, x$memory_global_by_current_bytes	Memory use, grouped by allocation type
memory_global_total, x$memory_global_total	Total memory use
metrics	Server metrics
processlist, x$processlist	Processlist information
ps_check_lost_instrumentation	Variables that have lost instruments
schema_auto_increment_columns	AUTO_INCREMENT column information
schema_index_statistics, x$schema_index_statistics	Index statistics
schema_object_overview	Types of objects within each schema
schema_redundant_indexes	Duplicate or redundant indexes
schema_table_lock_waits, x$schema_table_lock_waits	Sessions waiting for metadata locks
schema_table_statistics, x$schema_table_statistics	Table statistics
schema_table_statistics_with_buffer, x$schema_table_statistics_with_buffer	Table statistics, including InnoDB buffer pool statistics
schema_tables_with_full_table_scans, x$schema_tables_with_full_table_scans	Tables being accessed with full scans
schema_unused_indexes	Indexes not in active use
session, x$session	Processlist information for user sessions
session_ssl_status	Connection SSL information
statement_analysis, x$statement_analysis	Statement aggregate statistics
statements_with_errors_or_warnings, x$statements_with_errors_or_warnings	Statements that have produced errors or warnings
statements_with_full_table_scans, x$statements_with_full_table_scans	Statements that have done full table scans
statements_with_runtimes_in_95th_percentile, x$statements_with_runtimes_in_95th_percentile	Statements with highest average runtime

View Name	Description
statements_with_sorting, x$statements_with_sorting	Statements that performed sorts
statements_with_temp_tables, x$statements_with_temp_tables	Statements that used temporary tables
user_summary, x$user_summary	User statement and connection activity
user_summary_by_file_io, x$user_summary_by_file_io	File I/O, grouped by user
user_summary_by_file_io_type, x$user_summary_by_file_io_type	File I/O, grouped by user and event
user_summary_by_stages, x$user_summary_by_stages	Stage events, grouped by user
user_summary_by_statement_latency, x$user_summary_by_statement_latency	Statement statistics, grouped by user
user_summary_by_statement_type, x$user_summary_by_statement_type	Statements executed, grouped by user and statement
version	Current sys schema and MySQL server versions
wait_classes_global_by_avg_latency, x$wait_classes_global_by_avg_latency	Wait class average latency, grouped by event class
wait_classes_global_by_latency, x$wait_classes_global_by_latency	Wait class total latency, grouped by event class
waits_by_host_by_latency, x$waits_by_host_by_latency	Wait events, grouped by host and event
waits_by_user_by_latency, x$waits_by_user_by_latency	Wait events, grouped by user and event
waits_global_by_latency, x$waits_global_by_latency	Wait events, grouped by event
x$ps_digest_95th_percentile_by_avg_us	Helper view for 95th-percentile views
x$ps_digest_avg_latency_distribution	Helper view for 95th-percentile views
x$ps_schema_table_statistics_io	Helper view for table-statistics views
x$schema_flattened_keys	Helper view for schema_redundant_indexes

Table 26.3 sys Schema Stored Procedures

Procedure Name	Description
create_synonym_db()	Create synonym for schema
diagnostics()	Collect system diagnostic information
execute_prepared_stmt()	Execute prepared statement
ps_setup_disable_background_threads()	Disable background thread instrumentation
ps_setup_disable_consumer()	Disable consumers
ps_setup_disable_instrument()	Disable instruments
ps_setup_disable_thread()	Disable instrumentation for thread
ps_setup_enable_background_threads()	Enable background thread instrumentation
ps_setup_enable_consumer()	Enable consumers

Procedure Name	Description
ps_setup_enable_instrument()	Enable instruments
ps_setup_enable_thread()	Enable instrumentation for thread
ps_setup_reload_saved()	Reload saved Performance Schema configuration
ps_setup_reset_to_default()	Reset saved Performance Schema configuration
ps_setup_save()	Save Performance Schema configuration
ps_setup_show_disabled()	Display disabled Performance Schema configuration
ps_setup_show_disabled_consumers()	Display disabled Performance Schema consumers
ps_setup_show_disabled_instruments()	Display disabled Performance Schema instruments
ps_setup_show_enabled()	Display enabled Performance Schema configuration
ps_setup_show_enabled_consumers()	Display enabled Performance Schema consumers
ps_setup_show_enabled_instruments()	Display enabled Performance Schema instruments
ps_statement_avg_latency_histogram()	Display statement latency histogram
ps_trace_statement_digest()	Trace Performance Schema instrumentation for digest
ps_trace_thread()	Dump Performance Schema data for thread
ps_truncate_all_tables()	Truncate Performance Schema summary tables
statement_performance_analyzer()	Report of statements running on server
table_exists()	Whether a table exists

Table 26.4 sys Schema Stored Functions

Function Name	Description
extract_schema_from_file_name()	Extract schema name from file path name
extract_table_from_file_name()	Extract table name from file path name
format_bytes()	Convert byte value to value with units
format_path()	Replace data and temp-file directories in path name with symbolic values
format_statement()	Truncate long statement to fixed length
format_time()	Convert picoseconds value to value with units
list_add()	Add item to list
list_drop()	Remove item from list
ps_is_account_enabled()	Check whether account instrumentation is enabled
ps_is_consumer_enabled()	Check whether consumer is enabled

Function Name	Description
ps_is_instrument_default_enabled()	Check whether instrument is enabled
ps_is_instrument_default_timed()	Check whether instrument is timed
ps_is_thread_instrumented()	Check whether thread is instrumented
ps_thread_account()	Return account for thread ID
ps_thread_id()	Return thread ID for connection ID
ps_thread_stack()	Return event information for thread ID
ps_thread_trx_info()	Return transaction information for thread ID
quote_identifier()	Return string as quoted identifier
sys_get_config()	Return sys schema configuration option
version_major()	MySQL server major version number
version_minor()	MySQL server minor version number
version_patch()	MySQL server patch release version number

26.4.2 sys Schema Tables and Triggers

The following sections describe sys schema tables and triggers.

26.4.2.1 The sys_config Table

This table contains sys schema configuration options, one row per option. Configuration changes made by updating this table persist across client sessions and server restarts.

The sys_config table has these columns:

- variable

  The configuration option name.

- value

  The configuration option value.

- set_time

  The timestamp of the most recent modification to the row.

- set_by

  The account that made the most recent modification to the row. The value is NULL if the row has not been changed since the sys schema was installed.

As an efficiency measure to minimize the number of direct reads from the sys_config table, sys schema functions that use a value from this table check for a user-defined variable with a corresponding name, which is the user-defined variable having the same name plus a @sys. prefix. (For example, the variable corresponding to the diagnostics.include_raw option is @sys.diagnostics.include_raw.) If the user-defined variable exists in the current session and is non-NULL, the function uses its value in preference to the value in the sys_config table. Otherwise, the function reads and uses the value from the table. In the latter case, the calling function conventionally also sets the corresponding user-defined variable to the table value so that further references to the configuration option within the same session use the variable and need not read the table again.

For example, the `statement_truncate_len` option controls the maximum length of statements returned by the `format_statement()` function. The default is 64. To temporarily change the value to 32 for your current session, set the corresponding `@sys.statement_truncate_len` user-defined variable:

```
mysql> SET @stmt = 'SELECT variable, value, set_time, set_by FROM sys_config';
mysql> SELECT format_statement(@stmt);
+-------------------------------------------------------+
| format_statement(@stmt)                               |
+-------------------------------------------------------+
| SELECT variable, value, set_time, set_by FROM sys_config |
+-------------------------------------------------------+
mysql> SET @sys.statement_truncate_len = 32;
mysql> SELECT format_statement(@stmt);
+----------------------------------+
| format_statement(@stmt)          |
+----------------------------------+
| SELECT variabl ... ROM sys_config |
+----------------------------------+
```

Subsequent invocations of `format_statement()` within the session continue to use the user-defined variable value (32), rather than the value stored in the table (64).

To stop using the user-defined variable and revert to using the value in the table, set the variable to `NULL` within your session:

```
mysql> SET @sys.statement_truncate_len = NULL;
mysql> SELECT format_statement(@stmt);
+-------------------------------------------------------+
| format_statement(@stmt)                               |
+-------------------------------------------------------+
| SELECT variable, value, set_time, set_by FROM sys_config |
+-------------------------------------------------------+
```

Alternatively, end your current session (causing the user-defined variable to no longer exist) and begin a new session.

The conventional relationship just described between options in the `sys_config` table and user-defined variables can be exploited to make temporary configuration changes that end when your session ends. However, if you set a user-defined variable and then subsequently change the corresponding table value within the same session, the changed table value will not be used in that session as long as the user-defined variable exists with a non-`NULL` value. (The changed table value *will* be used in other sessions that do not have the user-defined variable assigned.)

The following list describes the options in the `sys_config` table and the corresponding user-defined variables:

- `diagnostics.allow_i_s_tables`, `@sys.diagnostics.allow_i_s_tables`

  If this option is `ON`, the `diagnostics()` procedure is permitted to perform table scans on the `INFORMATION_SCHEMA.TABLES` table. This can be expensive if there are many tables. The default is `OFF`.

- `diagnostics.include_raw`, `@sys.diagnostics.include_raw`

  If this option is `ON`, the `diagnostics()` procedure includes the raw output from querying the `metrics` view. The default is `OFF`.

- `ps_thread_trx_info.max_length`, `@sys.ps_thread_trx_info.max_length`

  The maximum length for JSON output produced by the `ps_thread_trx_info()` function. The default is 65535.

- `statement_performance_analyzer.limit`, `@sys.statement_performance_analyzer.limit`

  The maximum number of rows to return for views that have no built-in limit. (For example, the `statements_with_runtimes_in_95th_percentile` view has a built-in limit in the sense that it returns only statements with average execution time in the 95th percentile.) The default is 100.

- `statement_performance_analyzer.view`, `@sys.statement_performance_analyzer.view`

  The custom query or view to be used by the `statement_performance_analyzer()` procedure (which is itself invoked by the `diagnostics()` procedure). If the option value contains a space, it is interpreted as a query. Otherwise, it must be the name of an existing view that queries the Performance Schema `events_statements_summary_by_digest` table. There cannot be any `LIMIT` clause in the query or view definition if the `statement_performance_analyzer.limit` configuration option is greater than 0. The default is `NULL` (no custom view defined).

- `statement_truncate_len`, `@sys.statement_truncate_len`

  The maximum length of statements returned by the `format_statement()` function. Longer statements are truncated to this length. The default is 64.

Other options can be added to the `sys_config` table. For example, the `diagnostics()` and `execute_prepared_stmt()` procedures use the `debug` option if it exists, but this option is not part of the `sys_config` table by default because debug output normally is enabled only temporarily, by setting the corresponding `@sys.debug` user-defined variable. To enable debug output without having to set that variable in individual sessions, add the option to the table:

```
mysql> INSERT INTO sys_config (variable, value) VALUES('debug', 'ON');
```

To change the debug setting in the table, do two things. First, modify the value in the table itself:

```
mysql> UPDATE sys_config SET value = 'OFF' WHERE variable = 'debug';
```

Second, to also ensure that procedure invocations within the current session use the changed value from the table, set the corresponding user-defined variable to `NULL`:

```
mysql> SET @sys.debug = NULL;
```

26.4.2.2 The sys_config_insert_set_user Trigger

For rows added to the `sys_config` table by `INSERT` statements, the `sys_config_insert_set_user` trigger sets the `set_by` column to the current user.

26.4.2.3 The sys_config_update_set_user Trigger

The `sys_config_update_set_user` trigger for the `sys_config` table is similar to the `sys_config_insert_set_user` trigger, but for `UPDATE` statements.

26.4.3 sys Schema Views

The following sections describe `sys` schema views.

The `sys` schema contains many views that summarize Performance Schema tables in various ways. Most of these views come in pairs, such that one member of the pair has the same name as the other member, plus a `x$` prefix. For example, the `host_summary_by_file_io` view summarizes file I/O grouped by host and displays latencies converted from picoseconds to more readable values (with units);

```
```

```
mysql> SELECT * FROM host_summary_by_file_io;
+------------+-------+------------+
| host       | ios   | io_latency |
+------------+-------+------------+
| localhost  | 67570 | 5.38 s     |
| background |  3468 | 4.18 s     |
+------------+-------+------------+
```

The `x$host_summary_by_file_io` view summarizes the same data but displays unformatted picosecond latencies:

```
mysql> SELECT * FROM x$host_summary_by_file_io;
+------------+-------+---------------+
| host       | ios   | io_latency    |
+------------+-------+---------------+
| localhost  | 67574 | 5380678125144 |
| background |  3474 | 4758696829416 |
+------------+-------+---------------+
```

The view without the `x$` prefix is intended to provide output that is more user friendly and easier to read. The view with the `x$` prefix that displays the same values in raw form is intended more for use with other tools that perform their own processing on the data.

Views without the `x$` prefix differ from the corresponding `x$` views in these ways:

- Byte values are formatted with size units using `format_bytes()`.
- Time values are formatted with temporal units using `format_time()`.
- SQL statements are truncated to a maximum display width using `format_statement()`.
- Path name are shortened using `format_path()`.

26.4.3.1 The host_summary and x$host_summary Views

These views summarize statement activity, file I/O, and connections, grouped by host.

The `host_summary` and `x$host_summary` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `statements`

  The total number of statements for the host.

- `statement_latency`

  The total wait time of timed statements for the host.

- `statement_avg_latency`

  The average wait time per timed statement for the host.

- `table_scans`

  The total number of table scans for the host.

- `file_ios`

  The total number of file I/O events for the host.

- `file_io_latency`

  The total wait time of timed file I/O events for the host.

- `current_connections`

  The current number of connections for the host.

- `total_connections`

  The total number of connections for the host.

- `unique_users`

  The number of distinct users for the host.

- `current_memory`

  The current amount of allocated memory for the host.

- `total_memory_allocated`

  The total amount of allocated memory for the host.

26.4.3.2 The host_summary_by_file_io and x$host_summary_by_file_io Views

These views summarize file I/O, grouped by host. By default, rows are sorted by descending total file I/O latency.

The `host_summary_by_file_io` and `x$host_summary_by_file_io` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `ios`

  The total number of file I/O events for the host.

- `io_latency`

  The total wait time of timed file I/O events for the host.

26.4.3.3 The host_summary_by_file_io_type and x$host_summary_by_file_io_type Views

These views summarize file I/O, grouped by host and event type. By default, rows are sorted by host and descending total I/O latency.

The `host_summary_by_file_io_type` and `x$host_summary_by_file_io_type` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `event_name`

  The file I/O event name.

- `total`

  The total number of occurrences of the file I/O event for the host.

- `total_latency`

  The total wait time of timed occurrences of the file I/O event for the host.

- `max_latency`

  The maximum single wait time of timed occurrences of the file I/O event for the host.

26.4.3.4 The host_summary_by_stages and x$host_summary_by_stages Views

These views summarize statement stages, grouped by host. By default, rows are sorted by host and descending total latency.

The `host_summary_by_stages` and `x$host_summary_by_stages` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `event_name`

  The stage event name.

- `total`

  The total number of occurrences of the stage event for the host.

- `total_latency`

  The total wait time of timed occurrences of the stage event for the host.

- `avg_latency`

  The average wait time per timed occurrence of the stage event for the host.

26.4.3.5 The host_summary_by_statement_latency and x$host_summary_by_statement_latency Views

These views summarize overall statement statistics, grouped by host. By default, rows are sorted by descending total latency.

The `host_summary_by_statement_latency` and `x$host_summary_by_statement_latency` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `total`

  The total number of statements for the host.

- `total_latency`

  The total wait time of timed statements for the host.

- `max_latency`

  The maximum single wait time of timed statements for the host.

- `lock_latency`

  The total time waiting for locks by timed statements for the host.

- `rows_sent`

  The total number of rows returned by statements for the host.

- `rows_examined`

  The total number of rows read from storage engines by statements for the host.

- `rows_affected`

  The total number of rows affected by statements for the host.

- `full_scans`

  The total number of full table scans by statements for the host.

26.4.3.6 The host_summary_by_statement_type and x$host_summary_by_statement_type Views

These views summarize informaion about statements executed, grouped by host and statement type. By default, rows are sorted by host and descending total latency.

The `host_summary_by_statement_type` and `x$host_summary_by_statement_type` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `statement`

  The final component of the statement event name.

- `total`

  The total number of occurrences of the statement event for the host.

- `total_latency`

  The total wait time of timed occurrences of the statement event for the host.

- `max_latency`

  The maximum single wait time of timed occurrences of the statement event for the host.

- `lock_latency`

  The total time waiting for locks by timed occurrences of the statement event for the host.

- `rows_sent`

  The total number of rows returned by occurrences of the statement event for the host.

- `rows_examined`

The total number of rows read from storage engines by occurrences of the statement event for the host.

- `rows_affected`

  The total number of rows affected by occurrences of the statement event for the host.

- `full_scans`

  The total number of full table scans by occurrences of the statement event for the host.

26.4.3.7 The innodb_buffer_stats_by_schema and x$innodb_buffer_stats_by_schema Views

These views summarize the information in the `INFORMATION_SCHEMA INNODB_BUFFER_PAGE` table, grouped by schema. By default, rows are sorted by descending buffer size.

The `innodb_buffer_stats_by_schema` and `x$innodb_buffer_stats_by_schema` views have these columns:

- `object_schema`

  The schema name for the object, or `InnoDB System` if the table belongs to the `InnoDB` storage engine.

- `allocated`

  The total number of bytes allocated for the schema.

- `data`

  The total number of data bytes allocated for the schema.

- `pages`

  The total number of pages allocated for the schema.

- `pages_hashed`

  The total number of hashed pages allocated for the schema.

- `pages_old`

  The total number of old pages allocated for the schema.

- `rows_cached`

  The total number of cached rows for the schema.

26.4.3.8 The innodb_buffer_stats_by_table and x$innodb_buffer_stats_by_table Views

These views summarize the information in the `INFORMATION_SCHEMA INNODB_BUFFER_PAGE` table, grouped by schema and table. By default, rows are sorted by descending buffer size.

The `innodb_buffer_stats_by_table` and `x$innodb_buffer_stats_by_table` views have these columns:

- `object_schema`

  The schema name for the object, or `InnoDB System` if the table belongs to the `InnoDB` storage engine.

- `object_name`

  The table name.

- `allocated`

  The total number of bytes allocated for the table.

- `data`

  The number of data bytes allocated for the table.

- `pages`

  The total number of pages allocated for the table.

- `pages_hashed`

  The number of hashed pages allocated for the table.

- `pages_old`

  The number of old pages allocated for the table.

- `rows_cached`

  The number of cached rows for the table.

26.4.3.9 The innodb_lock_waits and x$innodb_lock_waits Views

These views summarize the `InnoDB` locks that transactions are waiting for. By default, rows are sorted by descending lock age.

The `innodb_lock_waits` and `x$innodb_lock_waits` views have these columns:

- `wait_started`

  The time at which the lock wait started.

- `wait_age`

  How long the lock has been waited for, as a `TIME` value.

- `wait_age_secs`

  How long the lock has been waited for, in seconds.

- `locked_table_schema`

  The schema that contains the locked table.

- `locked_table_name`

  The name of the locked table.

- `locked_table_partition`

  The name of the locked partition, if any; `NULL` otherwise.

- `locked_table_subpartition`

  The name of the locked subpartition, if any; `NULL` otherwise.

- `locked_index`

The name of the locked index.

- `locked_type`

  The type of the waiting lock.

- `waiting_trx_id`

  The ID of the waiting transaction.

- `waiting_trx_started`

  The time at which the waiting transaction started.

- `waiting_trx_age`

  How long the waiting transaction has been waiting, as a `TIME` value.

- `waiting_trx_rows_locked`

  The number of rows locked by the waiting transaction.

- `waiting_trx_rows_modified`

  The number of rows modified by the waiting transaction.

- `waiting_pid`

  The processlist ID of the waiting transaction.

- `waiting_query`

  The statement that is waiting for the lock.

- `waiting_lock_id`

  The ID of the waiting lock.

- `waiting_lock_mode`

  The mode of the waiting lock.

- `blocking_trx_id`

  The ID of the transaction that is blocking the waiting lock.

- `blocking_pid`

  The processlist ID of the blocking transaction.

- `blocking_query`

  The statement the blocking transaction is executing.

- `blocking_lock_id`

  The ID of the lock that is blocking the waiting lock.

- `blocking_lock_mode`

  The mode of the lock that is blocking the waiting lock.

- `blocking_trx_started`

The time at which the blocking transaction started.

- `blocking_trx_age`

  How long the blocking transaction has been executing, as a `TIME` value.

- `blocking_trx_rows_locked`

  The number of rows locked by the blocking transaction.

- `blocking_trx_rows_modified`

  The number of rows modified by the blocking transaction.

- `sql_kill_blocking_query`

  The `KILL` statement to execute to kill the blocking statement.

- `sql_kill_blocking_connection`

  The `KILL` statement to execute to kill the session running the blocking statement.

26.4.3.10 The io_by_thread_by_latency and x$io_by_thread_by_latency Views

These views summarize I/O consumers to display time waiting for I/O, grouped by thread. By default, rows are sorted by descending total I/O latency.

The `io_by_thread_by_latency` and `x$io_by_thread_by_latency` views have these columns:

- `user`

  For foreground threads, the account associated with the thread. For background threads, the thread name.

- `total`

  The total number of I/O events for the thread.

- `total_latency`

  The total wait time of timed I/O events for the thread.

- `min_latency`

  The minimum single wait time of timed I/O events for the thread.

- `avg_latency`

  The average wait time per timed I/O event for the thread.

- `max_latency`

  The maximum single wait time of timed I/O events for the thread.

- `thread_id`

  The thread ID.

- `processlist_id`

  For foreground threads, the processlist ID of the thread. For background threads, `NULL`.

26.4.3.11 The io_global_by_file_by_bytes and x$io_global_by_file_by_bytes Views

These views summarize global I/O consumers to display amount of I/O, grouped by file. By default, rows are sorted by descending total I/O (bytes read and written).

The `io_global_by_file_by_bytes` and `x$io_global_by_file_by_bytes` views have these columns:

- `file`

  The file path name.

- `count_read`

  The total number of read events for the file.

- `total_read`

  The total number of bytes read from the file.

- `avg_read`

  The average number of bytes per read from the file.

- `count_write`

  The total number of write events for the file.

- `total_written`

  The total number of bytes written to the file.

- `avg_write`

  The average number of bytes per write to the file.

- `total`

  The total number of bytes read and written for the file.

- `write_pct`

  The percentage of total bytes of I/O that were writes.

26.4.3.12 The io_global_by_file_by_latency and x$io_global_by_file_by_latency Views

These views summarize global I/O consumers to display time waiting for I/O, grouped by file. By default, rows are sorted by descending total latency.

The `io_global_by_file_by_latency` and `x$io_global_by_file_by_latency` views have these columns:

- `file`

  The file path name.

- `total`

  The total number of I/O events for the file.

- `total_latency`

  The total wait time of timed I/O events for the file.

- `count_read`

The total number of read I/O events for the file.

- `read_latency`

  The total wait time of timed read I/O events for the file.

- `count_write`

  The total number of write I/O events for the file.

- `write_latency`

  The total wait time of timed write I/O events for the file.

- `count_misc`

  The total number of other I/O events for the file.

- `misc_latency`

  The total wait time of timed other I/O events for the file.

26.4.3.13 The io_global_by_wait_by_bytes and x$io_global_by_wait_by_bytes Views

These views summarize global I/O consumers to display amount of I/O and time waiting for I/O, grouped by event. By default, rows are sorted by descending total I/O (bytes read and written).

The `io_global_by_wait_by_bytes` and `x$io_global_by_wait_by_bytes` views have these columns:

- `event_name`

  The I/O event name, with the `wait/io/file/` prefix stripped.

- `total`

  The total number of occurrences of the I/O event.

- `total_latency`

  The total wait time of timed occurrences of the I/O event.

- `min_latency`

  The minimum single wait time of timed occurrences of the I/O event.

- `avg_latency`

  The average wait time per timed occurrence of the I/O event.

- `max_latency`

  The maximum single wait time of timed occurrences of the I/O event.

- `count_read`

  The number of read requests for the I/O event.

- `total_read`

  The number of bytes read for the I/O event.

- `avg_read`

The average number of bytes per read for the I/O event.

- `count_write`

  The number of write requests for the I/O event.

- `total_written`

  The number of bytes written for the I/O event.

- `avg_written`

  The average number of bytes per write for the I/O event.

- `total_requested`

  The total number of bytes read and written for the I/O event.

26.4.3.14 The io_global_by_wait_by_latency and x$io_global_by_wait_by_latency Views

These views summarize global I/O consumers to display amount of I/O and time waiting for I/O, grouped by event. By default, rows are sorted by descending total latency.

The `io_global_by_wait_by_latency` and `x$io_global_by_wait_by_latency` views have these columns:

- `event_name`

  The I/O event name, with the `wait/io/file/` prefix stripped.

- `total`

  The total number of occurrences of the I/O event.

- `total_latency`

  The total wait time of timed occurrences of the I/O event.

- `avg_latency`

  The average wait time per timed occurrence of the I/O event.

- `max_latency`

  The maximum single wait time of timed occurrences of the I/O event.

- `read_latency`

  The total wait time of timed read occurrences of the I/O event.

- `write_latency`

  The total wait time of timed write occurrences of the I/O event.

- `misc_latency`

  The total wait time of timed other occurrences of the I/O event.

- `count_read`

  The number of read requests for the I/O event.

- `total_read`

The number of bytes read for the I/O event.

- `avg_read`

  The average number of bytes per read for the I/O event.

- `count_write`

  The number of write requests for the I/O event.

- `total_written`

  The number of bytes written for the I/O event.

- `avg_written`

  The average number of bytes per write for the I/O event.

26.4.3.15 The latest_file_io and x$latest_file_io Views

These views summarize file I/O activity, grouped by file and thread. By default, rows are sorted with most recent I/O first.

The `latest_file_io` and `x$latest_file_io` views have these columns:

- `thread`

  For foreground threads, the account associated with the thread (and port number for TCP/IP connections). For background threads, the thread name and thread ID

- `file`

  The file path name.

- `latency`

  The wait time of the file I/O event.

- `operation`

  The type of operation.

- `requested`

  The number of data bytes requested for the file I/O event.

26.4.3.16 The memory_by_host_by_current_bytes and x$memory_by_host_by_current_bytes Views

These views summarize memory use, grouped by host. By default, rows are sorted by descending amount of memory used.

The `memory_by_host_by_current_bytes` and `x$memory_by_host_by_current_bytes` views have these columns:

- `host`

  The host from which the client connected. Rows for which the `HOST` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `current_count_used`

The current number of allocated memory blocks that have not been freed yet for the host.

- `current_allocated`

  The current number of allocated bytes that have not been freed yet for the host.

- `current_avg_alloc`

  The current number of allocated bytes per memory block for the host.

- `current_max_alloc`

  The largest single current memory allocation in bytes for the host.

- `total_allocated`

  The total memory allocation in bytes for the host.

26.4.3.17 The memory_by_thread_by_current_bytes and x$memory_by_thread_by_current_bytes Views

These views summarize memory use, grouped by thread. By default, rows are sorted by descending amount of memory used.

The `memory_by_thread_by_current_bytes` and `x$memory_by_thread_by_current_bytes` views have these columns:

- `thread_id`

  The thread ID.

- `user`

  The thread user or thread name.

- `current_count_used`

  The current number of allocated memory blocks that have not been freed yet for the thread.

- `current_allocated`

  The current number of allocated bytes that have not been freed yet for the thread.

- `current_avg_alloc`

  The current number of allocated bytes per memory block for the thread.

- `current_max_alloc`

  The largest single current memory allocation in bytes for the thread.

- `total_allocated`

  The total memory allocation in bytes for the thread.

26.4.3.18 The memory_by_user_by_current_bytes and x$memory_by_user_by_current_bytes Views

These views summarize memory use, grouped by user. By default, rows are sorted by descending amount of memory used.

The `memory_by_user_by_current_bytes` and `x$memory_by_user_by_current_bytes` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `current_count_used`

  The current number of allocated memory blocks that have not been freed yet for the user.

- `current_allocated`

  The current number of allocated bytes that have not been freed yet for the user.

- `current_avg_alloc`

  The current number of allocated bytes per memory block for the user.

- `current_max_alloc`

  The largest single current memory allocation in bytes for the user.

- `total_allocated`

  The total memory allocation in bytes for the user.

26.4.3.19 The memory_global_by_current_bytes and x$memory_global_by_current_bytes Views

These views summarize memory use, grouped by allocation type (that is, by event). By default, rows are sorted by descending amount of memory used.

The `memory_global_by_current_bytes` and `x$memory_global_by_current_bytes` views have these columns:

- `event_name`

  The memory event name.

- `current_count`

  The total number of occurrences of the event.

- `current_alloc`

  The current number of allocated bytes that have not been freed yet for the event.

- `current_avg_alloc`

  The current number of allocated bytes per memory block for the event.

- `high_count`

  The high-water mark for number of memory blocks allocated for the event.

- `high_alloc`

  The high-water mark for number of bytes allocated for the event.

- `high_avg_alloc`

  The high-water mark for average number of bytes per memory block allocated for the event.

26.4.3.20 The memory_global_total and x$memory_global_total Views

These views summarize total memory use within the server.

The `memory_global_total` and `x$memory_global_total` views have these columns:

- `total_allocated`

  The total bytes of memory allocated within the server.

26.4.3.21 The metrics View

This view summarizes MySQL server metrics to show variable names, values, types, and whether they are enabled. By default, rows are sorted by variable type and name.

The `metrics` view includes this information:

- Global status variables from the Performance Schema `global_status` table
- `InnoDB` metrics from the `INFORMATION_SCHEMA INNODB_METRICS` table
- Current and total memory allocation, based on the Performance Schema memory instrumentation
- The current time (human readable and Unix timestamp formats)

There is some duplication of information between the `global_status` and `INNODB_METRICS` tables, which the `metrics` view eliminates.

The `metrics` view has these columns:

- `Variable_name`

  The metric name. The metric type determines the source from which the name is taken:

  - For global status variables: The `VARIABLE_NAME` column of the `global_status` table
  - For `InnoDB` metrics: The `NAME` column of the `INNODB_METRICS` table
  - For other metrics: A view-provided descriptive string

- `Variable_value`

  The metric value. The metric type determines the source from which the value is taken:

  - For global status variables: The `VARIABLE_VALUE` column of the `global_status` table
  - For `InnoDB` metrics: The `COUNT` column of the `INNODB_METRICS` table
  - For memory metrics: The relevant column from the Performance Schema `memory_summary_global_by_event_name` table
  - For the current time: The value of `NOW(3)` or `UNIX_TIMESTAMP(NOW(3))`

- `Type`

  The metric type:

  - For global status variables: `Global Status`
  - For `InnoDB` metrics: `InnoDB Metrics - %`, where `%` is replaced by the value of the `SUBSYSTEM` column of the `INNODB_METRICS` table
  - For memory metrics: `Performance Schema`

  - For the current time: `System Time`

- `Enabled`

  Whether the metric is enabled:

  - For global status variables: `YES`
  - For `InnoDB` metrics: `YES` if the `STATUS` column of the `INNODB_METRICS` table is `enabled`, `NO` otherwise
  - For memory metrics: `NO`, `YES`, or `PARTIAL` (currently, `PARTIAL` occurs only for memory metrics and indicates that not all `memory/%` instruments are enabled; Performance Schema memory instruments are always enabled)
  - For the current time: `YES`

26.4.3.22 The processlist and x$processlist Views

These views summarize processlist information. They provide more complete information than the `SHOW PROCESSLIST` statement and the `INFORMATION_SCHEMA PROCESSLIST` table, and are also nonblocking. By default, rows are sorted by descending process time and descending wait time.

The column descriptions here are brief. For additional information, see the description of the Performance Schema `threads` table at Section 25.11.16.3, “The threads Table”.

The `processlist` and `x$processlist` views have these columns:

- `thd_id`

  The thread ID.

- `conn_id`

  The connection ID.

- `user`

  The thread user or thread name.

- `db`

  The default database for the thread, or `NULL` if there is none.

- `command`

  For foreground threads, the type of command the thread is executing on behalf of the client, or `Sleep` if the session is idle.

- `state`

  An action, event, or state that indicates what the thread is doing.

- `time`

  The time in seconds that the thread has been in its current state.

- `current_statement`

  The statement the thread is executing, or `NULL` if it is not executing any statement.

- `statement_latency`

How long the statement has been executing.

- `progress`

  The percentage of work completed for stages that support progress reporting. See Section 26.3, “sys Schema Progress Reporting”.

- `lock_latency`

  The time spent waiting for locks by the current statement.

- `rows_examined`

  The number of rows read from storage engines by the current statement.

- `rows_sent`

  The number of rows returned by the current statement.

- `rows_affected`

  The number of rows affected by the current statement.

- `tmp_tables`

  The number of internal in-memory temporary tables created by the current statement.

- `tmp_disk_tables`

  The number of internal on-disk temporary tables created by the current statement.

- `full_scan`

  The number of full table scans performed by the current statement.

- `last_statement`

  The last statement executed by the thread, if there is no currently executing statement or wait.

- `last_statement_latency`

  How long the last statement executed.

- `current_memory`

  The number of bytes allocated by the thread.

- `last_wait`

  The name of the most recent wait event for the thread.

- `last_wait_latency`

  The wait time of the most recent wait event for the thread.

- `source`

  The source file and line number containing the instrumented code that produced the event.

- `trx_latency`

  The wait time of the current transaction for the thread.

- `trx_state`

The state for the current transaction for the thread.

- `trx_autocommit`

  Whether autocommit mode was enabled when the current transaction started.

- `pid`

  The client process ID.

- `program_name`

  The client program name.

26.4.3.23 The ps_check_lost_instrumentation View

This view returns information about lost Performance Schema instruments, to indicate whether the Performance Schema is unable to monitor all runtime data.

The `ps_check_lost_instrumentation` view has these columns:

- `variable_name`

  The Performance Schema status variable name indicating which type of instrument was lost.

- `variable_value`

  The number of instruments lost.

26.4.3.24 The schema_auto_increment_columns View

This view indicates which tables have `AUTO_INCREMENT` columns and provides information about those columns, such as the current and maximum column values and the usage ratio (ratio of used to possible values). By default, rows are sorted by descending usage ratio and maximum column value.

Tables in these schemas are excluded from view output: `mysql`, `sys`, `INFORMATION_SCHEMA`, `performance_schema`.

The `schema_auto_increment_columns` view has these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table that contains the `AUTO_INCREMENT` column.

- `column_name`

  The name of the `AUTO_INCREMENT` column.

- `data_type`

  The data type of the column.

- `column_type`

  The column type of the column, which is the data type plus possibly other information. For example, for a column with a `bigint(20) unsigned` column type, the data type is just `bigint`.

- `is_signed`

Whether the column type is signed.

- `is_unsigned`

  Whether the column type is unsigned.

- `max_value`

  The maximum permitted value for the column.

- `auto_increment`

  The current `AUTO_INCREMENT` value for the column.

- `auto_increment_ratio`

  The ratio of used to permitted values for the column. This indicates how much of the sequence of values is “used up.”

26.4.3.25 The schema_index_statistics and x$schema_index_statistics Views

These views provide index statistics. By default, rows are sorted by descending total index latency.

The `schema_index_statistics` and `x$schema_index_statistics` views have these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table that contains the index.

- `index_name`

  The name of the index.

- `rows_selected`

  The total number of rows read using the index.

- `select_latency`

  The total wait time of timed reads using the index.

- `rows_inserted`

  The total number of rows inserted into the index.

- `insert_latency`

  The total wait time of timed inserts into the index.

- `rows_updated`

  The total number of rows updated in the index.

- `update_latency`

  The total wait time of timed updates in the index.

- `rows_deleted`

The total number of rows deleted from the index.

- `delete_latency`

  The total wait time of timed deletes from the index.

26.4.3.26 The schema_object_overview View

This view summarizes the types of objects within each schema. By default, rows are sorted by schema and object type.

Note

For MySQL instances with a large number of objects, this view might take a long time to execute.

The `schema_object_overview` view has these columns:

- `db`

  The schema name.

- `object_type`

  The object type: `BASE TABLE`, `INDEX (index_type)`, `EVENT`, `FUNCTION`, `PROCEDURE`, `TRIGGER`, `VIEW`.

- `count`

  The number of objects in the schema of the given type.

26.4.3.27 The schema_redundant_indexes and x$schema_flattened_keys Views

The `schema_redundant_indexes` view displays indexes that duplicate other indexes or are made redundant by them. The `x$schema_flattened_keys` view is a helper view for `schema_redundant_indexes`.

In the following column descriptions, the dominant index is the one that makes the redundant index redundant.

The `schema_redundant_indexes` view has these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table that contains the index.

- `redundant_index_name`

  The name of the redundant index.

- `redundant_index_columns`

  The names of the columns in the redundant index.

- `redundant_index_non_unique`

  The number of nonunique columns in the redundant index.

- `dominant_index_name`

  The name of the dominant index.

- `dominant_index_columns`

  The names of the columns in the dominant index.

- `dominant_index_non_unique`

  The number of nonunique columns in the dominant index.

- `subpart_exists`

  Whether the index indexes only part of a column.

- `sql_drop_index`

  The statement to execute to drop the redundant index.

The `x$schema_flattened_keys` view has these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table that contains the index.

- `index_name`

  An index name.

- `non_unique`

  The number of nonunique columns in the index.

- `subpart_exists`

  Whether the index indexes only part of a column.

- `index_columns`

  The name of the columns in the index.

26.4.3.28 The schema_table_lock_waits and x$schema_table_lock_waits Views

These views display which sessions are blocked waiting on metadata locks, and what is blocking them.

The column descriptions here are brief. For additional information, see the description of the Performance Schema `metadata_locks` table at Section 25.11.12.3, “The metadata_locks Table”.

The `schema_table_lock_waits` and `x$schema_table_lock_waits` views have these columns:

- `object_schema`

  The schema containing the object to be locked.

- `object_name`

  The name of the instrumented object.

- `waiting_thread_id`

The thread ID of the thread that is waiting for the lock.

- `waiting_pid`

  The processlist ID of the thread that is waiting for the lock.

- `waiting_account`

  The account associated with the session that is waiting for the lock.

- `waiting_lock_type`

  The type of the waiting lock.

- `waiting_lock_duration`

  How long the waiting lock has been waiting.

- `waiting_query`

  The statement that is waiting for the lock.

- `waiting_query_secs`

  How long the statement has been waiting, in seconds.

- `waiting_query_rows_affected`

  The number of rows affected by the statement.

- `waiting_query_rows_examined`

  The number of rows read from storage engines by the statement.

- `blocking_thread_id`

  The thread ID of the thread that is blocking the waiting lock.

- `blocking_pid`

  The processlist ID of the thread that is blocking the waiting lock.

- `blocking_account`

  The account associated with the thread that is blocking the waiting lock.

- `blocking_lock_type`

  The type of lock that is blocking the waiting lock.

- `blocking_lock_duration`

  How long the blocking lock has been held.

- `sql_kill_blocking_query`

  The `KILL` statement to execute to kill the blocking statement.

- `sql_kill_blocking_connection`

  The `KILL` statement to execute to kill the session running the blocking statement.

26.4.3.29 The schema_table_statistics and x$schema_table_statistics Views

These views summarize table statistics. By default, rows are sorted by descending total wait time (tables with most contention first).

These views user a helper view, `x$ps_schema_table_statistics_io`.

The `schema_table_statistics` and `x$schema_table_statistics` views have these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table name.

- `total_latency`

  The total wait time of timed I/O events for the table.

- `rows_fetched`

  The total number of rows read from the table.

- `fetch_latency`

  The total wait time of timed read I/O events for the table.

- `rows_inserted`

  The total number of rows inserted into the table.

- `insert_latency`

  The total wait time of timed insert I/O events for the table.

- `rows_updated`

  The total number of rows updated in the table.

- `update_latency`

  The total wait time of timed update I/O events for the table.

- `rows_deleted`

  The total number of rows deleted from the table.

- `delete_latency`

  The total wait time of timed delete I/O events for the table.

- `io_read_requests`

  The total number of read requests for the table.

- `io_read`

  The total number of bytes read from the table.

- `io_read_latency`

  The total wait time of reads from the table.

- `io_write_requests`

The total number of write requests for the table.

- `io_write`

  The total number of bytes written to the table.

- `io_write_latency`

  The total wait time of writes to the table.

- `io_misc_requests`

  The total number of miscellaneous I/O requests for the table.

- `io_misc_latency`

  The total wait time of miscellaneous I/O requests for the table.

26.4.3.30 The schema_table_statistics_with_buffer and x $schema_table_statistics_with_buffer Views

These views summarize table statistics, including `InnoDB` buffer pool statistics. By default, rows are sorted by descending total wait time (tables with most contention first).

These views user a helper view, `x$ps_schema_table_statistics_io`.

The `schema_table_statistics_with_buffer` and `x $schema_table_statistics_with_buffer` views have these columns:

- `table_schema`

  The schema that contains the table.

- `table_name`

  The table name.

- `rows_fetched`

  The total number of rows read from the table.

- `fetch_latency`

  The total wait time of timed read I/O events for the table.

- `rows_inserted`

  The total number of rows inserted into the table.

- `insert_latency`

  The total wait time of timed insert I/O events for the table.

- `rows_updated`

  The total number of rows updated in the table.

- `update_latency`

  The total wait time of timed update I/O events for the table.

- `rows_deleted`

The total number of rows deleted from the table.

- `delete_latency`

  The total wait time of timed delete I/O events for the table.

- `io_read_requests`

  The total number of read requests for the table.

- `io_read`

  The total number of bytes read from the table.

- `io_read_latency`

  The total wait time of reads from the table.

- `io_write_requests`

  The total number of write requests for the table.

- `io_write`

  The total number of bytes written to the table.

- `io_write_latency`

  The total wait time of writes to the table.

- `io_misc_requests`

  The total number of miscellaneous I/O requests for the table.

- `io_misc_latency`

  The total wait time of miscellaneous I/O requests for the table.

- `innodb_buffer_allocated`

  The total number of `InnoDB` buffer bytes allocated for the table.

- `innodb_buffer_data`

  The total number of `InnoDB` data bytes allocated for the table.

- `innodb_buffer_free`

  The total number of `InnoDB` nondata bytes allocated for the table (`innodb_buffer_allocated` − `innodb_buffer_data`).

- `innodb_buffer_pages`

  The total number of `InnoDB` pages allocated for the table.

- `innodb_buffer_pages_hashed`

  The total number of `InnoDB` hashed pages allocated for the table.

- `innodb_buffer_pages_old`

  The total number of `InnoDB` old pages allocated for the table.

- `innodb_buffer_rows_cached`

The total number of `InnoDB` cached rows for the table.

26.4.3.31 The schema_tables_with_full_table_scans and x$schema_tables_with_full_table_scans Views

These views display which tables are being accessed with full table scans. By default, rows are sorted by descending rows scanned.

The `schema_tables_with_full_table_scans` and `x$schema_tables_with_full_table_scans` views have these columns:

- `object_schema`

  The schema name.

- `object_name`

  The table name.

- `rows_full_scanned`

  The total number of rows scanned by full scans of the table.

- `latency`

  The total wait time of full scans of the table.

26.4.3.32 The schema_unused_indexes View

These views display indexes for which there are no events, which indicates that they are not being used. By default, rows are sorted by schema and table.

This view is most useful when the server has been up and processing long enough that its workload is representative. Otherwise, presence of an index in this view may not be meaningful.

The `schema_unused_indexes` view has these columns:

- `object_schema`

  The schema name.

- `object_name`

  The table name.

- `index_name`

  The unused index name.

26.4.3.33 The session and x$session Views

These views are similar to `processlist` and `x$processlist`, but they filter out background processes to display only user sessions. For descriptions of the columns, see Section 26.4.3.22, “The processlist and x$processlist Views”.

26.4.3.34 The session_ssl_status View

For each connection, this view displays the SSL version, cipher, and count of reused SSL sessions.

The `session_ssl_status` view has these columns:

- `thread_id`

The thread ID for the connection.

- `ssl_version`

  The version of SSL used for the connection.

- `ssl_cipher`

  The SSL cipher used for the connection.

- `ssl_sessions_reused`

  The number of reused SSL sessions for the connection.

26.4.3.35 The statement_analysis and x$statement_analysis Views

These views list normalized statements with aggregated statistics. The content mimics the MySQL Enterprise Monitor Query Analysis view. By default, rows are sorted by descending total latency.

The `statement_analysis` and `x$statement_analysis` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `full_scan`

  The total number of full table scans performed by occurrences of the statement.

- `exec_count`

  The total number of times the statement has executed.

- `err_count`

  The total number of errors produced by occurrences of the statement.

- `warn_count`

  The total number of warnings produced by occurrences of the statement.

- `total_latency`

  The total wait time of timed occurrences of the statement.

- `max_latency`

  The maximum single wait time of timed occurrences of the statement.

- `avg_latency`

  The average wait time per timed occurrence of the statement.

- `lock_latency`

  The total time waiting for locks by timed occurrences of the statement.

- `rows_sent`

The total number of rows returned by occurrences of the statement.

- `rows_sent_avg`

  The average number of rows returned per occurrence of the statement.

- `rows_examined`

  The total number of rows read from storage engines by occurrences of the statement.

- `rows_examined_avg`

  The average number of rows read from storage engines per occurrence of the statement.

- `rows_affected`

  The total number of rows affected by occurrences of the statement.

- `rows_affected_avg`

  The average number of rows affected per occurrence of the statement.

- `tmp_tables`

  The total number of internal in-memory temporary tables created by occurrences of the statement.

- `tmp_disk_tables`

  The total number of internal on-disk temporary tables created by occurrences of the statement.

- `rows_sorted`

  The total number of rows sorted by occurrences of the statement.

- `sort_merge_passes`

  The total number of sort merge passes by occurrences of the statement.

- `digest`

  The statement digest.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

26.4.3.36 The statements_with_errors_or_warnings and x$statements_with_errors_or_warnings Views

These views display normalized statements that have produced errors or warnings. By default, rows are sorted by descending error and warning counts.

The `statements_with_errors_or_warnings` and `x$statements_with_errors_or_warnings` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `exec_count`

  The total number of times the statement has executed.

- `errors`

  The total number of errors produced by occurrences of the statement.

- `error_pct`

  The percentage of statement occurrences that produced errors.

- `warnings`

  The total number of warnings produced by occurrences of the statement.

- `warning_pct`

  The percentage of statement occurrences that produced warnings.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

- `digest`

  The statement digest.

26.4.3.37 The statements_with_full_table_scans and x$statements_with_full_table_scans Views

These views display normalized statements that have done full table scans. By default, rows are sorted by descending percentage of time a full scan was done and descending total latency.

The `statements_with_full_table_scans` and `x$statements_with_full_table_scans` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `exec_count`

  The total number of times the statement has executed.

- `total_latency`

  The total wait time of timed statement events for the statement.

- `no_index_used_count`

  The total number of times no index was used to scan the table.

- `no_good_index_used_count`

  The total number of times no good index was used to scan the table.

- `no_index_used_pct`

  The percentage of the time no index was used to scan the table.

- `rows_sent`

  The total number of rows returned from the table.

- `rows_examined`

  The total number of rows read from the storage engine for the table.

- `rows_sent_avg`

  The average number of rows returned from the table.

- `rows_examined_avg`

  The average number of rows read from the storage engine for the table.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

- `digest`

  The statement digest.

26.4.3.38 The statements_with_runtimes_in_95th_percentile and x$statements_with_runtimes_in_95th_percentile Views

These views list statements with runtimes in the 95th percentile. By default, rows are sorted by descending average latency.

Both views use two helper views, `x$ps_digest_avg_latency_distribution` and `x$ps_digest_95th_percentile_by_avg_us`.

The `statements_with_runtimes_in_95th_percentile` and `x$statements_with_runtimes_in_95th_percentile` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `full_scan`

  The total number of full table scans performed by occurrences of the statement.

- `exec_count`

  The total number of times the statement has executed.

- `err_count`

  The total number of errors produced by occurrences of the statement.

- `warn_count`

  The total number of warnings produced by occurrences of the statement.

- `total_latency`

  The total wait time of timed occurrences of the statement.

- `max_latency`

  The maximum single wait time of timed occurrences of the statement.

- `avg_latency`

  The average wait time per timed occurrence of the statement.

- `rows_sent`

  The total number of rows returned by occurrences of the statement.

- `rows_sent_avg`

  The average number of rows returned per occurrence of the statement.

- `rows_examined`

  The total number of rows read from storage engines by occurrences of the statement.

- `rows_examined_avg`

  The average number of rows read from storage engines per occurrence of the statement.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

- `digest`

  The statement digest.

26.4.3.39 The statements_with_sorting and x$statements_with_sorting Views

These views list normalized statements that have performed sorts. By default, rows are sorted by descending total latency.

The `statements_with_sorting` and `x$statements_with_sorting` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `exec_count`

The total number of times the statement has executed.

- `total_latency`

  The total wait time of timed occurrences of the statement.

- `sort_merge_passes`

  The total number of sort merge passes by occurrences of the statement.

- `avg_sort_merges`

  The average number of sort merge passes per occurrence of the statement.

- `sorts_using_scans`

  The total number of sorts using table scans by occurrences of the statement.

- `sort_using_range`

  The total number of sorts using range accesses by occurrences of the statement.

- `rows_sorted`

  The total number of rows sorted by occurrences of the statement.

- `avg_rows_sorted`

  The average number of rows sorted per occurrence of the statement.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

- `digest`

  The statement digest.

26.4.3.40 The statements_with_temp_tables and x$statements_with_temp_tables Views

These views list normalized statements that have used temporary tables. By default, rows are sorted by descending number of on-disk temporary tables used and descending number of in-memory temporary tables used.

The `statements_with_temp_tables` and `x$statements_with_temp_tables` views have these columns:

- `query`

  The normalized statement string.

- `db`

  The default database for the statement, or `NULL` if there is none.

- `exec_count`

  The total number of times the statement has executed.

- `total_latency`

  The total wait time of timed occurrences of the statement.

- `memory_tmp_tables`

  The total number of internal in-memory temporary tables created by occurrences of the statement.

- `disk_tmp_tables`

  The total number of internal on-disk temporary tables created by occurrences of the statement.

- `avg_tmp_tables_per_query`

  The average number of internal temporary tables created per occurrence of the statement.

- `tmp_tables_to_disk_pct`

  The percentage of internal in-memory temporary tables that were converted to on-disk tables.

- `first_seen`

  The time at which the statement was first seen.

- `last_seen`

  The time at which the statement was most recently seen.

- `digest`

  The statement digest.

26.4.3.41 The user_summary and x$user_summary Views

These views summarize statement activity, file I/O, and connections, grouped by user. By default, rows are sorted by descending total latency.

The `user_summary` and `x$user_summary` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `statements`

  The total number of statements for the user.

- `statement_latency`

  The total wait time of timed statements for the user.

- `statement_avg_latency`

  The average wait time per timed statement for the user.

- `table_scans`

  The total number of table scans for the user.

- `file_ios`

  The total number of file I/O events for the user.

- `file_io_latency`

  The total wait time of timed file I/O events for the user.

- `current_connections`

  The current number of connections for the user.

- `total_connections`

  The total number of connections for the user.

- `unique_hosts`

  The number of distinct hosts from which connections for the user have originated.

- `current_memory`

  The current amount of allocated memory for the user.

- `total_memory_allocated`

  The total amount of allocated memory for the user.

26.4.3.42 The user_summary_by_file_io and x$user_summary_by_file_io Views

These views summarize file I/O, grouped by user. By default, rows are sorted by descending total file I/O latency.

The `user_summary_by_file_io` and `x$user_summary_by_file_io` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `ios`

  The total number of file I/O events for the user.

- `io_latency`

  The total wait time of timed file I/O events for the user.

26.4.3.43 The user_summary_by_file_io_type and x$user_summary_by_file_io_type Views

These views summarize file I/O, grouped by user and event type. By default, rows are sorted by user and descending total latency.

The `user_summary_by_file_io_type` and `x$user_summary_by_file_io_type` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `event_name`

  The file I/O event name.

- `total`

  The total number of occurrences of the file I/O event for the user.

- `latency`

  The total wait time of timed occurrences of the file I/O event for the user.

- `max_latency`

  The maximum single wait time of timed occurrences of the file I/O event for the user.

26.4.3.44 The user_summary_by_stages and x$user_summary_by_stages Views

These views summarize stages, grouped by user. By default, rows are sorted by user and descending total stage latency.

The `user_summary_by_stages` and `x$user_summary_by_stages` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `event_name`

  The stage event name.

- `total`

  The total number of occurrences of the stage event for the user.

- `total_latency`

  The total wait time of timed occurrences of the stage event for the user.

- `avg_latency`

  The average wait time per timed occurrence of the stage event for the user.

26.4.3.45 The user_summary_by_statement_latency and x$user_summary_by_statement_latency Views

These views summarize overall statement statistics, grouped by user. By default, rows are sorted by descending total latency.

The `user_summary_by_statement_latency` and `x$user_summary_by_statement_latency` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `total`

  The total number of statements for the user.

- `total_latency`

  The total wait time of timed statements for the user.

- `max_latency`

  The maximum single wait time of timed statements for the user.

- `lock_latency`

  The total time waiting for locks by timed statements for the user.

- `rows_sent`

  The total number of rows returned by statements for the user.

- `rows_examined`

  The total number of rows read from storage engines by statements for the user.

- `rows_affected`

  The total number of rows affected by statements for the user.

- `full_scans`

  The total number of full table scans by statements for the user.

26.4.3.46 The user_summary_by_statement_type and x$user_summary_by_statement_type Views

These views summarize informaion about statements executed, grouped by user and statement type. By default, rows are sorted by user and descending total latency.

The `user_summary_by_statement_type` and `x$user_summary_by_statement_type` views have these columns:

- `user`

  The client user name. Rows for which the `USER` column in the underlying Performance Schema table is `NULL` are assumed to be for background threads and are reported with a host name of `background`.

- `statement`

  The final component of the statement event name.

- `total`

  The total number of occurrences of the statement event for the user.

- `total_latency`

  The total wait time of timed occurrences of the statement event for the user.

- `max_latency`

  The maximum single wait time of timed occurrences of the statement event for the user.

- `lock_latency`

  The total time waiting for locks by timed occurrences of the statement event for the user.

- `rows_sent`

  The total number of rows returned by occurrences of the statement event for the user.

- `rows_examined`

The total number of rows read from storage engines by occurrences of the statement event for the user.

- `rows_affected`

  The total number of rows affected by occurrences of the statement event for the user.

- `full_scans`

  The total number of full table scans by occurrences of the statement event for the user.

26.4.3.47 The version View

This view provides the current `sys` schema and MySQL server versions.

The `version` view has these columns:

- `sys_version`

  The `sys` schema version.

- `mysql_version`

  The MySQL server version.

26.4.3.48 The wait_classes_global_by_avg_latency and x$wait_classes_global_by_avg_latency Views

These views summarize wait class average latencies, grouped by event class. By default, rows are sorted by descending average latency. Idle events are ignored.

An event class is determined by stripping from the event name everything after the first three components. For example, the class for `wait/io/file/sql/slow_log` is `wait/io/file`.

The `wait_classes_global_by_avg_latency` and `x$wait_classes_global_by_avg_latency` views have these columns:

- `event_class`

  The event class.

- `total`

  The total number of occurrences of events in the class.

- `total_latency`

  The total wait time of timed occurrences of events in the class.

- `min_latency`

  The minimum single wait time of timed occurrences of events in the class.

- `avg_latency`

  The average wait time per timed occurrence of events in the class.

- `max_latency`

  The maximum single wait time of timed occurrences of events in the class.

26.4.3.49 The wait_classes_global_by_latency and x$wait_classes_global_by_latency Views

These views summarize wait class total latencies, grouped by event class. By default, rows are sorted by descending total latency. Idle events are ignored.

An event class is determined by stripping from the event name everything after the first three components. For example, the class for `wait/io/file/sql/slow_log` is `wait/io/file`.

The `wait_classes_global_by_latency` and `x$wait_classes_global_by_latency` views have these columns:

- `event_class`

  The event class.

- `total`

  The total number of occurrences of events in the class.

- `total_latency`

  The total wait time of timed occurrences of events in the class.

- `min_latency`

  The minimum single wait time of timed occurrences of events in the class.

- `avg_latency`

  The average wait time per timed occurrence of events in the class.

- `max_latency`

  The maximum single wait time of timed occurrences of events in the class.

26.4.3.50 The waits_by_host_by_latency and x$waits_by_host_by_latency Views

These views summarize wait events, grouped by host and event. By default, rows are sorted by host and descending total latency. Idle events are ignored.

The `waits_by_host_by_latency` and `x$waits_by_host_by_latency` views have these columns:

- `host`

  The host from which the connection originated.

- `event`

  The event name.

- `total`

  The total number of occurrences of the event for the host.

- `total_latency`

  The total wait time of timed occurrences of the event for the host.

- `avg_latency`

  The average wait time per timed occurrence of the event for the host.

- `max_latency`

  The maximum single wait time of timed occurrences of the event for the host.

26.4.3.51 The waits_by_user_by_latency and x$waits_by_user_by_latency Views

These views summarize wait events, grouped by user and event. By default, rows are sorted by user and descending total latency. Idle events are ignored.

The `waits_by_user_by_latency` and `x$waits_by_user_by_latency` views have these columns:

- `user`

  The user associated with the connection.

- `event`

  The event name.

- `total`

  The total number of occurrences of the event for the user.

- `total_latency`

  The total wait time of timed occurrences of the event for the user.

- `avg_latency`

  The average wait time per timed occurrence of the event for the user.

- `max_latency`

  The maximum single wait time of timed occurrences of the event for the user.

26.4.3.52 The waits_global_by_latency and x$waits_global_by_latency Views

These views summarize wait events, grouped by event. By default, rows are sorted by descending total latency. Idle events are ignored.

The `waits_global_by_latency` and `x$waits_global_by_latency` views have these columns:

- `events`

  The event name.

- `total`

  The total number of occurrences of the event.

- `total_latency`

  The total wait time of timed occurrences of the event.

- `avg_latency`

  The average wait time per timed occurrence of the event.

- `max_latency`

  The maximum single wait time of timed occurrences of the event.

26.4.4 sys Schema Stored Procedures

The following sections describe `sys` schema stored procedures.

26.4.4.1 The create_synonym_db() Procedure

Given a schema name, this procedure creates a synonym schema containing views that refer to all the tables and views in the original schema. This can be used, for example, to create a shorter name by which to refer to a schema with a long name (such as `info` rather than `INFORMATION_SCHEMA`).

Parameters

- `in_db_name VARCHAR(64)`: The name of the schema for which to create the synonym.
- `in_synonym VARCHAR(64)`: The name to use for the synonym schema. This schema must not already exist.

Example

```
mysql> SHOW DATABASES;
+--------------------+
| Database           |
+--------------------+
| information_schema |
| mysql              |
| performance_schema |
| sys                |
| world              |
+--------------------+
mysql> CALL create_synonym_db('INFORMATION_SCHEMA', 'info');
+--------------------------------------+
| summary                              |
+--------------------------------------+
| Created 63 views in the info database |
+--------------------------------------+
mysql> SHOW DATABASES;
+--------------------+
| Database           |
+--------------------+
| information_schema |
| info               |
| mysql              |
| performance_schema |
| sys                |
| world              |
+--------------------+
mysql> SHOW FULL TABLES FROM info;
+---------------------------------------+------------+
| Tables_in_info                        | Table_type |
+---------------------------------------+------------+
| character_sets                        | VIEW       |
| collation_character_set_applicability | VIEW       |
| collations                            | VIEW       |
| column_privileges                     | VIEW       |
| columns                               | VIEW       |
...
```

26.4.4.2 The diagnostics() Procedure

Creates a report of the current server status for diagnostic purposes.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Data collected for `diagnostics()` includes this information:

- Information from the `metrics` view (see Section 26.4.3.21, “The metrics View”)
- Information from other relevant `sys` schema views, such as the one that detemines queries in the 95th percentile

- Information from the `ndbinfo` schema, if the MySQL server is part of MySQL Cluster
- Replication status (both master and slave)

Some of the sys schema views are calculated as initial (optional), overall, and delta values:

- The initial view is the content of the view at the start of the `diagnostics()` procedure. This output is the same as the start values used for the delta view. The initial view is included if the `diagnostics.include_raw` configuration option is `ON`.
- The overall view is the content of the view at the end of the `diagnostics()` procedure. This output is the same as the end values used for the delta view. The overall view is always included.
- The delta view is the difference from the beginning to the end of procedure execution. The minimum and maximum values are the minimum and maximum values from the end view, respectively. They do not necessarily reflect the minimum and maximum values in the monitored period. Except for the `metrics` view, the delta is calculated only between the first and last outputs.

Parameters

- `in_max_runtime INT UNSIGNED`: The maximum data collection time in seconds. Use `NULL` to collect data for the default of 60 seconds. Otherwise, use a value greater than 0.
- `in_interval INT UNSIGNED`: The sleep time between data collections in seconds. Use `NULL` to sleep for the default of 30 seconds. Otherwise, use a value greater than 0.
- `in_auto_config ENUM('current', 'medium', 'full')`: The Performance Schema configuration to use. Permitted values are:
  - `current`: Use the current instrument and consumer settings.
  - `medium`: Enable some instruments and consumers.
  - `full`: Enable all instruments and consumers.

Note

The more instruments and consumers enabled, the more impact on MySQL server performance. Be careful with the `medium` setting and especially the `full` setting, which has a large performance impact.

Use of the `medium` or `full` setting requires the `SUPER` privilege.

If a setting other than `current` is chosen, the current settings are restored at the end of the procedure.

Configuration Options

`diagnostics()` operation can be modified using the following configuration options or their corresponding user-defined variables (see Section 26.4.2.1, “The sys_config Table”):

- `debug`, `@sys.debug`

  If this option is `ON`, produce debugging output. The default is `OFF`.

- `diagnostics.allow_i_s_tables`, `@sys.diagnostics.allow_i_s_tables`

  If this option is `ON`, the `diagnostics()` procedure is permitted to perform table scans on the `INFORMATION_SCHEMA.TABLES` table. This can be expensive if there are many tables. The default is `OFF`.

- `diagnostics.include_raw`, `@sys.diagnostics.include_raw`

If this option is `ON`, the `diagnostics()` procedure output includes the raw output from querying the `metrics` view. The default is `OFF`.

- `statement_truncate_len`, `@sys.statement_truncate_len`

  The maximum length of statements returned by the `format_statement()` function. Longer statements are truncated to this length. The default is 64.

Example

Create a diagnostics report that starts an iteration every 30 seconds and runs for at most 120 seconds using the current Performance Schema settings:

```
mysql> CALL diagnostics(120, 30, 'current');
```

To capture the output from the `diagnostics()` procedure in a file as it runs, use the `mysql` client `tee filename` and `notee` commands (see Section 4.5.1.2, “mysql Commands”):

```
mysql> tee diag.out;
mysql> CALL diagnostics(120, 30, 'current');
mysql> notee;
```

26.4.4.3 The execute_prepared_stmt() Procedure

Given an SQL statement as a string, executes it as a prepared statement. The prepared statement is deallocated after execution, so it is not subject to reuse. Thus, this procedure is useful primarily for executing dynamic statements on a one-time basis.

This procedure uses `sys_execute_prepared_stmt` as the prepared statement name. If that statement name exists when the procedure is called, its previous content is destroyed.

Parameters

- `in_query LONGTEXT CHARACTER SET utf8`: The statement string to execute.

Configuration Options

`execute_prepared_stmt()` operation can be modified using the following configuration options or their corresponding user-defined variables (see Section 26.4.2.1, “The sys_config Table”):

- `debug`, `@sys.debug`

  If this option is `ON`, produce debugging output. The default is `OFF`.

Example

```
mysql> CALL execute_prepared_stmt('SELECT COUNT(*) FROM mysql.user');
+----------+
| COUNT(*) |
+----------+
|       15 |
+----------+
```

26.4.4.4 The ps_setup_disable_background_threads() Procedure

Disables Performance Schema instrumentation for all background threads. Produces a result set indicating how many background threads were disabled. Already disabled threads do not count.

Parameters

None.

Example

```
mysql> CALL ps_setup_disable_background_threads();
+------------------------------+
| summary                      |
+------------------------------+
| Disabled 24 background threads |
+------------------------------+
```

26.4.4.5 The ps_setup_disable_consumer() Procedure

Disables Performance Schema consumers with names that contain the argument. Produces a result set indicating how many consumers were disabled. Already disabled consumers do not count.

Parameters

- `consumer VARCHAR(128)`: The value used to match consumer names, which are identified by using `%consumer%` as an operand for a `LIKE` pattern match.

  A value of `''` matches all consumers.

Example

Disable all statement consumers:

```
mysql> CALL ps_setup_disable_consumer('statement');
+----------------------+
| summary              |
+----------------------+
| Disabled 4 consumers |
+----------------------+
```

26.4.4.6 The ps_setup_disable_instrument() Procedure

Disables Performance Schema instruments with names that contain the argument. Produces a result set indicating how many instruments were disabled. Already disabled instruments do not count.

Parameters

- `in_pattern VARCHAR(128)`: The value used to match instrument names, which are identified by using `%in_pattern%` as an operand for a `LIKE` pattern match.

  A value of `''` matches all instruments.

Example

Disable a specific instrument:

```
mysql> CALL ps_setup_disable_instrument('wait/lock/metadata/sql/mdl');
+----------------------+
| summary              |
+----------------------+
| Disabled 1 instrument |
+----------------------+
```

Disable all mutex instruments:

```
mysql> CALL ps_setup_disable_instrument('mutex');
+--------------------------+
| summary                  |
+--------------------------+
| Disabled 177 instruments |
```

```
+-------------------------+
```

26.4.4.7 The ps_setup_disable_thread() Procedure

Given a connection ID, disables Performance Schema instrumentation for the thread. Produces a result set indicating how many threads were disabled. Already disabled threads do not count.

Parameters

- `in_connection_id BIGINT`: The connection ID. This is a connection ID as given in the `PROCESSLIST_ID` column of the Performance Schema `threads` table or the `Id` column of `SHOW PROCESSLIST` output.

Example

Disable a specific connection by its connection ID:

```
mysql> CALL ps_setup_disable_thread(225);
+-------------------+
| summary           |
+-------------------+
| Disabled 1 thread |
+-------------------+
```

Disable the current connection:

```
mysql> CALL ps_setup_disable_thread(CONNECTION_ID());
+-------------------+
| summary           |
+-------------------+
| Disabled 1 thread |
+-------------------+
```

26.4.4.8 The ps_setup_enable_background_threads() Procedure

Enables Performance Schema instrumentation for all background threads. Produces a result set indicating how many background threads were enabled. Already enabled threads do not count.

Parameters

None.

Example

```
mysql> CALL ps_setup_enable_background_threads();
+------------------------------+
| summary                      |
+------------------------------+
| Enabled 24 background threads |
+------------------------------+
```

26.4.4.9 The ps_setup_enable_consumer() Procedure

Enables Performance Schema consumers with names that contain the argument. Produces a result set indicating how many consumers were enabled. Already enabled consumers do not count.

Parameters

- `consumer VARCHAR(128)`: The value used to match consumer names, which are identified by using `%consumer%` as an operand for a `LIKE` pattern match.

  A value of `''` matches all consumers.

Example

Enable all statement consumers:

```
mysql> CALL ps_setup_enable_consumer('statement');
+--------------------+
| summary            |
+--------------------+
| Enabled 4 consumers |
+--------------------+
```

26.4.4.10 The ps_setup_enable_instrument() Procedure

Enables Performance Schema instruments with names that contain the argument. Produces a result set indicating how many instruments were enabled. Already enabled instruments do not count.

Parameters

- `in_pattern VARCHAR(128)`: The value used to match instrument names, which are identified by using `%in_pattern%` as an operand for a `LIKE` pattern match.

  A value of `''` matches all instruments.

Example

Enable a specific instrument:

```
mysql> CALL ps_setup_enable_instrument('wait/lock/metadata/sql/mdl');
+---------------------+
| summary             |
+---------------------+
| Enabled 1 instrument |
+---------------------+
```

Enable all mutex instruments:

```
mysql> CALL ps_setup_enable_instrument('mutex');
+-------------------------+
| summary                 |
+-------------------------+
| Enabled 177 instruments |
+-------------------------+
```

26.4.4.11 The ps_setup_enable_thread() Procedure

Given a connection ID, enables Performance Schema instrumentation for the thread. Produces a result set indicating how many threads were enabled. Already enabled threads do not count.

Parameters

- `in_connection_id BIGINT`: The connection ID. This is a connection ID as given in the `PROCESSLIST_ID` column of the Performance Schema `threads` table or the `Id` column of `SHOW PROCESSLIST` output.

Example

Enable a specific connection by its connection ID:

```
mysql> CALL ps_setup_enable_thread(225);
+-----------------+
| summary         |
+-----------------+
| Enabled 1 thread |
```

```
+------------------+
```

Enable the current connection:

```
mysql> CALL ps_setup_enable_thread(CONNECTION_ID());
+------------------+
| summary          |
+------------------+
| Enabled 1 thread |
+------------------+
```

26.4.4.12 The ps_setup_reload_saved() Procedure

Reloads a Performance Schema configuration saved earlier within the same session using `ps_setup_save()`. For more information, see the description of `ps_setup_save()`.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Parameters

None.

26.4.4.13 The ps_setup_reset_to_default() Procedure

Resets the Performance Schema configuration to its default settings.

Parameters

- `in_verbose BOOLEAN`: Whether to display information about each setup stage during procedure execution. This includes the SQL statements executed.

Example

```
mysql> CALL ps_setup_reset_to_default(TRUE)\G
*************************** 1. row ***************************
status: Resetting: setup_actors
DELETE
FROM performance_schema.setup_actors
WHERE NOT (HOST = '%' AND USER = '%' AND ROLE = '%')

*************************** 1. row ***************************
status: Resetting: setup_actors
INSERT IGNORE INTO performance_schema.setup_actors
VALUES ('%', '%', '%')

...
```

26.4.4.14 The ps_setup_save() Procedure

Saves the current Performance Schema configuration. This enables you to alter the configuration temporarily for debugging or other purposes, then restore it to the previous state by invoking the `ps_setup_reload_saved()` procedure.

To prevent other simultaneous calls to save the configuration, `ps_setup_save()` acquires an advisory lock named `sys.ps_setup_save` by calling the `GET_LOCK()` function. `ps_setup_save()` takes a timeout parameter to indicate how many seconds to wait if the lock already exists (which indicates that some other session has a saved configuration outstanding). If the timeout expires without obtaining the lock, `ps_setup_save()` fails.

It is intended you call `ps_setup_reload_saved()` later within the *same* session as `ps_setup_save()` because the configuration is saved in `TEMPORARY` tables. `ps_setup_save()` drops the temporary tables and releases the lock. If you end your session without invoking `ps_setup_save()`, the tables and lock disappear automatically.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Parameters

- `in_timeout INT`: How many seconds to wait to obtain the `sys.ps_setup_save` lock. A negative timeout value means infinite timeout.

Example

```
mysql> CALL ps_setup_save(10);

... make Performance Schema configuration changes ...

mysql> CALL ps_setup_reload_saved();
```

26.4.4.15 The ps_setup_show_disabled() Procedure

Displays all currently disabled Performance Schema configuration.

Parameters

- `in_show_instruments BOOLEAN`: Whether to display disabled instruments. This might be a long list.
- `in_show_threads BOOLEAN`: Whether to display disabled threads.

Example

```
mysql> CALL ps_setup_show_disabled(TRUE, TRUE);
+----------------------------+
| performance_schema_enabled |
+----------------------------+
|                          1 |
+----------------------------+

+---------------+
| enabled_users |
+---------------+
| '%'@'%'       |
+---------------+

+-------------+----------------------+---------+-------+
| object_type | objects              | enabled | timed |
+-------------+----------------------+---------+-------+
| EVENT       | mysql.%              | NO      | NO    |
| EVENT       | performance_schema.% | NO      | NO    |
| EVENT       | information_schema.% | NO      | NO    |
| FUNCTION    | mysql.%              | NO      | NO    |
| FUNCTION    | performance_schema.% | NO      | NO    |
| FUNCTION    | information_schema.% | NO      | NO    |
| PROCEDURE   | mysql.%              | NO      | NO    |
| PROCEDURE   | performance_schema.% | NO      | NO    |
| PROCEDURE   | information_schema.% | NO      | NO    |
| TABLE       | mysql.%              | NO      | NO    |
| TABLE       | performance_schema.% | NO      | NO    |
| TABLE       | information_schema.% | NO      | NO    |
| TRIGGER     | mysql.%              | NO      | NO    |
| TRIGGER     | performance_schema.% | NO      | NO    |
| TRIGGER     | information_schema.% | NO      | NO    |
+-------------+----------------------+---------+-------+

...
```

26.4.4.16 The ps_setup_show_disabled_consumers() Procedure

Displays all currently disabled Performance Schema consumers.

Parameters

None.

Example

```
mysql> CALL ps_setup_show_disabled_consumers();
+---------------------------------+
| disabled_consumers              |
+---------------------------------+
| events_stages_current           |
| events_stages_history           |
| events_stages_history_long      |
| events_statements_history       |
| events_statements_history_long  |
| events_transactions_history     |
| events_transactions_history_long |
| events_waits_current            |
| events_waits_history            |
| events_waits_history_long       |
+---------------------------------+
```

26.4.4.17 The ps_setup_show_disabled_instruments() Procedure

Displays all currently disabled Performance Schema instruments. This might be a long list.

Parameters

None.

Example

```
mysql> CALL ps_setup_show_disabled_instruments()\G
*************************** 1. row ***************************
disabled_instruments: wait/synch/mutex/sql/TC_LOG_MMAP::LOCK_tc
               timed: NO
*************************** 2. row ***************************
disabled_instruments: wait/synch/mutex/sql/LOCK_des_key_file
               timed: NO
*************************** 3. row ***************************
disabled_instruments: wait/synch/mutex/sql/MYSQL_BIN_LOG::LOCK_commit
               timed: NO
...
```

26.4.4.18 The ps_setup_show_enabled() Procedure

Displays all currently enabled Performance Schema configuration.

Parameters

- `in_show_instruments BOOLEAN`: Whether to display enabled instruments. This might be a long list.
- `in_show_threads BOOLEAN`: Whether to display enabled threads.

Example

```
mysql> CALL ps_setup_show_enabled(FALSE, FALSE);
+---------------------------+
| performance_schema_enabled |
+---------------------------+
|                         1 |
+---------------------------+
1 row in set (0.00 sec)
```

```
+---------------+
| enabled_users |
+---------------+
| '%'@'%'       |
+---------------+
1 row in set (0.00 sec)

+-------------+---------------------+---------+-------+
| object_type | objects             | enabled | timed |
+-------------+---------------------+---------+-------+
| EVENT       | %.%                 | YES     | YES   |
| FUNCTION    | %.%                 | YES     | YES   |
| PROCEDURE   | %.%                 | YES     | YES   |
| TABLE       | %.%                 | YES     | YES   |
| TRIGGER     | %.%                 | YES     | YES   |
+-------------+---------------------+---------+-------+
5 rows in set (0.01 sec)

+----------------------------+
| enabled_consumers          |
+----------------------------+
| events_statements_current  |
| events_transactions_current |
| global_instrumentation     |
| thread_instrumentation     |
| statements_digest          |
+----------------------------+
```

26.4.4.19 The ps_setup_show_enabled_consumers() Procedure

Displays all currently enabled Performance Schema consumers.

Parameters

None.

Example

```
mysql> CALL ps_setup_show_enabled_consumers();
+----------------------------+
| enabled_consumers          |
+----------------------------+
| events_statements_current  |
| events_transactions_current |
| global_instrumentation     |
| thread_instrumentation     |
| statements_digest          |
+----------------------------+
```

26.4.4.20 The ps_setup_show_enabled_instruments() Procedure

Displays all currently enabled Performance Schema instruments. This might be a long list.

Parameters

None.

Example

```
mysql> CALL ps_setup_show_enabled_instruments()\G
*************************** 1. row ***************************
enabled_instruments: wait/io/file/sql/map
              timed: YES
*************************** 2. row ***************************
enabled_instruments: wait/io/file/sql/binlog
              timed: YES
*************************** 3. row ***************************
enabled_instruments: wait/io/file/sql/binlog_cache
```

```
              timed: YES
...
```

26.4.4.21 The ps_statement_avg_latency_histogram() Procedure

Displays a textual histogram graph of the average latency values across all normalized statements tracked within the Performance Schema `events_statements_summary_by_digest` table.

This procedure can be used to display a very high-level picture of the latency distribution of statements running within this MySQL instance.

Parameters

None.

Example

The histogram output in statement units. For example, `* = 2 units` in the histogram legend means that each `*` character represents 2 statements.

```
mysql> CALL ps_statement_avg_latency_histogram()\G
*************************** 1. row ***************************
Performance Schema Statement Digest Average Latency Histogram:

  . = 1 unit
  * = 2 units
  # = 3 units

(0 - 66ms)     88  | ##############################
(66 - 133ms)   14  | ..............
(133 - 199ms)  4   | ....
(199 - 265ms)  5   | **
(265 - 332ms)  1   | .
(332 - 398ms)  0   |
(398 - 464ms)  1   | .
(464 - 531ms)  0   |
(531 - 597ms)  0   |
(597 - 663ms)  0   |
(663 - 730ms)  0   |
(730 - 796ms)  0   |
(796 - 863ms)  0   |
(863 - 929ms)  0   |
(929 - 995ms)  0   |
(995 - 1062ms) 0   |

  Total Statements: 114; Buckets: 16; Bucket Size: 66 ms;
```

26.4.4.22 The ps_trace_statement_digest() Procedure

Traces all Performance Schema instrumentation for a specific statement digest.

If you find a statement of interest within the Performance Schema `events_statements_summary_by_digest` table, specify its `DIGEST` column MD5 value to this procedure and indicate the polling duration and interval. The result is a report of all statistics tracked within Performance Schema for that digest for the interval.

The procedure also attempts to execute `EXPLAIN` for the longest running example of the digest during the interval. This attempt might fail because the Performance Schema truncates long `SQL_TEXT` values. Consequently, `EXPLAIN` will fail due to parse errors.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Parameters

- `in_digest VARCHAR(32)`: The statement digest identifier to analyze.

- `in_runtime INT`: How long to run the analysis in seconds.
- `in_interval DECIMAL(2,2)`: The analysis interval in seconds (which can be fractional) at which to try to take snapshots.
- `in_start_fresh BOOLEAN`: Whether to truncate the Performance Schema `events_statements_history_long` and `events_stages_history_long` tables before starting.
- `in_auto_enable BOOLEAN`: Whether to automatically enable required consumers.

Example

```
mysql> CALL ps_trace_statement_digest('891ec6860f98ba46d89dd20b0c03652c', 10, 0.1, TRUE, TRUE);
+--------------------+
| SUMMARY STATISTICS |
+--------------------+
| SUMMARY STATISTICS |
+--------------------+
1 row in set (9.11 sec)

+------------+-----------+-----------+-----------+---------------+------------+------------+
| executions | exec_time | lock_time | rows_sent | rows_examined | tmp_tables | full_scans |
+------------+-----------+-----------+-----------+---------------+------------+------------+
|         21 | 4.11 ms   | 2.00 ms   |         0 |            21 |          0 |          0 |
+------------+-----------+-----------+-----------+---------------+------------+------------+
1 row in set (9.11 sec)

+----------------------------------------+-------+-----------+
| event_name                             | count | latency   |
+----------------------------------------+-------+-----------+
| stage/sql/checking query cache for query |    16 | 724.37 us |
| stage/sql/statistics                   |    16 | 546.92 us |
| stage/sql/freeing items                |    18 | 520.11 us |
| stage/sql/init                         |    51 | 466.80 us |
...
| stage/sql/cleaning up                  |    18 | 11.92 us  |
| stage/sql/executing                    |    16 | 6.95 us   |
+----------------------------------------+-------+-----------+
17 rows in set (9.12 sec)

+---------------------------+
| LONGEST RUNNING STATEMENT |
+---------------------------+
| LONGEST RUNNING STATEMENT |
+---------------------------+
1 row in set (9.16 sec)

+-----------+-----------+-----------+-----------+---------------+------------+-----------+
| thread_id | exec_time | lock_time | rows_sent | rows_examined | tmp_tables | full_scan |
+-----------+-----------+-----------+-----------+---------------+------------+-----------+
|    166646 | 618.43 us | 1.00 ms   |         0 |             1 |          0 |         0 |
+-----------+-----------+-----------+-----------+---------------+------------+-----------+
1 row in set (9.16 sec)

# Truncated for clarity...
+--------------------------------------------------------------------+
| sql_text                                                           |
+--------------------------------------------------------------------+
| select hibeventhe0_.id as id1382_, hibeventhe0_.createdTime ... |
+--------------------------------------------------------------------+
1 row in set (9.17 sec)

+----------------------------------------+-----------+
| event_name                             | latency   |
+----------------------------------------+-----------+
| stage/sql/init                         | 8.61 us   |
| stage/sql/Waiting for query cache lock | 453.23 us |
| stage/sql/init                         | 331.07 ns |
| stage/sql/checking query cache for query | 43.04 us  |
```

```
...
| stage/sql/freeing items                  | 30.46 us  |
| stage/sql/cleaning up                    | 662.13 ns |
+------------------------------------------+-----------+
18 rows in set (9.23 sec)

+----+-------------+--------------+-------+---------------+-----------+---------+-------------+------+-----
| id | select_type | table        | type  | possible_keys | key       | key_len | ref         | rows | Extr
+----+-------------+--------------+-------+---------------+-----------+---------+-------------+------+-----
|  1 | SIMPLE      | hibeventhe0_ | const | fixedTime     | fixedTime | 775     | const,const |    1 | NULL
+----+-------------+--------------+-------+---------------+-----------+---------+-------------+------+-----
1 row in set (9.27 sec)

Query OK, 0 rows affected (9.28 sec)
```

26.4.4.23 The ps_trace_thread() Procedure

Dumps all Performance Schema data for an instrumented thread to a `.dot` formatted graph file (for the DOT graph description language). Each result set returned from the procedure should be used for a complete graph.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Parameters

- `in_thread_id INT`: The thread to trace.
- `in_outfile VARCHAR(255)`: The name to use for the `.dot` output file.
- `in_max_runtime DECIMAL(20,2)`: The maximum number of seconds (which can be fractional) to collect data. Use `NULL` to collect data for the default of 60 seconds.
- `in_interval DECIMAL(20,2)`: The number of seconds (which can be fractional) to sleep between data collections. Use `NULL` to sleep for the default of 1 second.
- `in_start_fresh BOOLEAN`: Whether to reset all Performance Schema data before tracing.
- `in_auto_setup BOOLEAN`: Whether to disable all other threads and enable all instruments and consumers. This also resets the settings at the end of the run.
- `in_debug BOOLEAN`: Whether to include `file:lineno` information in the graph.

Example

```
mysql> CALL ps_trace_thread(25, CONCAT('/tmp/stack-', REPLACE(NOW(), ' ', '-'), '.dot'), NULL, NULL, TRUE,
+--------------------+
| summary            |
+--------------------+
| Disabled 1 thread  |
+--------------------+
1 row in set (0.00 sec)

+-------------------------------------------+
| Info                                      |
+-------------------------------------------+
| Data collection starting for THREAD_ID = 25 |
+-------------------------------------------+
1 row in set (0.03 sec)

+-----------------------------------------------------------+
| Info                                                      |
+-----------------------------------------------------------+
| Stack trace written to /tmp/stack-2014-02-16-21:18:41.dot |
+-----------------------------------------------------------+
1 row in set (60.07 sec)
```

```
+-----------------------------------------------------------------+
| Convert to PDF                                                  |
+-----------------------------------------------------------------+
| dot -Tpdf -o /tmp/stack_25.pdf /tmp/stack-2014-02-16-21:18:41.dot |
+-----------------------------------------------------------------+
1 row in set (60.07 sec)

+-----------------------------------------------------------------+
| Convert to PNG                                                  |
+-----------------------------------------------------------------+
| dot -Tpng -o /tmp/stack_25.png /tmp/stack-2014-02-16-21:18:41.dot |
+-----------------------------------------------------------------+
1 row in set (60.07 sec)

+------------------+
| summary          |
+------------------+
| Enabled 1 thread |
+------------------+
1 row in set (60.32 sec)
```

26.4.4.24 The ps_truncate_all_tables() Procedure

Truncates all Performance Schema summary tables, resetting all aggregated instrumentation as a snapshot. Produces a result set indicating how many tables were truncated.

Parameters

- `in_verbose BOOLEAN`: Whether to display each `TRUNCATE TABLE` statement before executing it.

Example

```
mysql> CALL ps_truncate_all_tables(FALSE);
+--------------------+
| summary            |
+--------------------+
| Truncated 44 tables |
+--------------------+
```

26.4.4.25 The statement_performance_analyzer() Procedure

Creates a report of the statements running on the server. The views are calculated based on the overall and/or delta activity.

This procedure requires the `SYSTEM_VARIABLES_ADMIN` or `SUPER` privilege because it manipulates the session `sql_log_bin` system variable to disable binary logging during its execution.

Parameters

- `in_action ENUM('snapshot', 'overall', 'delta', 'create_tmp', 'create_table', 'save', 'cleanup')`: The action to take. These values are permitted:
  - `snapshot`: Store a snapshot. The default is to make a snapshot of the current content of the Performance Schema `events_statements_summary_by_digest` table. By setting `in_table`, this can be overwritten to copy the content of the specified table. The snapshot is stored in the `sys` schema `tmp_digests` temporary table.
  - `overall`: Generate an analysis based on the content of the table specified by `in_table`. For the overall analysis, `in_table` can be `NOW()` to use a fresh snapshot. This overwrites an existing snapshot. Use `NULL` for `in_table` to use the existing snapshot. If `in_table` is `NULL` and no snapshot exists, a new snapshot is created. The `in_views` parameter and the `statement_performance_analyzer.limit` configuration option affect the operation of this procedure.

  - `delta`: Generate a delta analysis. The delta is calculated between the reference table specified by `in_table` and the snapshot, which must exist. This action uses the `sys` schema `tmp_digests_delta` temporary table. The `in_views` parameter and the `statement_performance_analyzer.limit` configuration option affect the operation of this procedure.
  - `create_table`: Create a regular table suitable for storing the snapshot for later use (for example, for calculating deltas).
  - `create_tmp`: Create a temporary table suitable for storing the snapshot for later use (for example, for calculating deltas).
  - `save`: Save the snapshot in the table specified by `in_table`. The table must exist and have the correct structure. If no snapshot exists, a new snapshot is created.
  - `cleanup`: Remove the temporary tables used for the snapshot and delta.

- `in_table VARCHAR(129)`: The table parameter used for some of the actions specified by the `in_action` parameter. Use the format *`db_name.tbl_name`* or *`tbl_name`* without using any backtick (`` ` ``) identifier-quoting characters. Periods (`.`) are not supported in database and table names.

  The meaning of the `in_table` value for each `in_action` value is detailed in the individual `in_action` value descriptions.

- `in_views SET ('with_runtimes_in_95th_percentile', 'analysis', 'with_errors_or_warnings', 'with_full_table_scans', 'with_sorting', 'with_temp_tables', 'custom')`: Which views to include. This parameter is a `SET` value, so it can contain multiple view names, separated by commas. The default is to include all views except `custom`. The following values are permitted:

  - `with_runtimes_in_95th_percentile`: Use the `statements_with_runtimes_in_95th_percentile` view.
  - `analysis`: Use the `statement_analysis` view.
  - `with_errors_or_warnings`: Use the `statements_with_errors_or_warnings` view.
  - `with_full_table_scans`: Use the `statements_with_full_table_scans` view.
  - `with_sorting`: Use the `statements_with_sorting` view.
  - `with_temp_tables`: Use the `statements_with_temp_tables` view.
  - `custom`: Use a custom view. This view must be specified using the `statement_performance_analyzer.view` configuration option to name a query or an existing view.

Configuration Options

`statement_performance_analyzer()` operation can be modified using the following configuration options or their corresponding user-defined variables (see Section 26.4.2.1, “The sys_config Table”):

- `debug`, `@sys.debug`

  If this option is `ON`, produce debugging output. The default is `OFF`.

- `statement_performance_analyzer.limit`, `@sys.statement_performance_analyzer.limit`

  The maximum number of rows to return for views that have no built-in limit. The default is 100.

- `statement_performance_analyzer.view`, `@sys.statement_performance_analyzer.view`

  The custom query or view to be used. If the option value contains a space, it is interpreted as a query. Otherwise, it must be the name of an existing view that queries the Performance Schema `events_statements_summary_by_digest` table. There cannot be any `LIMIT` clause in the query or view definition if the `statement_performance_analyzer.limit` configuration option is greater than 0. If specifying a view, use the same format as for the `in_table` parameter. The default is `NULL` (no custom view defined).

Example

To create a report with the queries in the 95th percentile since the last truncation of `events_statements_summary_by_digest` and with a one-minute delta period:

1. Create a temporary table to store the initial snapshot.
2. Create the initial snapshot.
3. Save the initial snapshot in the temporary table.
4. Wait one minute.
5. Create a new snapshot.
6. Perform analysis based on the new snapshot.
7. Perform analysis based on the delta between the initial and new snapshots.

```
mysql> CALL statement_performance_analyzer('create_tmp', 'mydb.tmp_digests_ini', NULL);
Query OK, 0 rows affected (0.08 sec)

mysql> CALL statement_performance_analyzer('snapshot', NULL, NULL);
Query OK, 0 rows affected (0.02 sec)

mysql> CALL statement_performance_analyzer('save', 'mydb.tmp_digests_ini', NULL);
Query OK, 0 rows affected (0.00 sec)

mysql> DO SLEEP(60);
Query OK, 0 rows affected (1 min 0.00 sec)

mysql> CALL statement_performance_analyzer('snapshot', NULL, NULL);
Query OK, 0 rows affected (0.02 sec)

mysql> CALL statement_performance_analyzer('overall', NULL, 'with_runtimes_in_95th_percentile');
+----------------------------------------+
| Next Output                            |
+----------------------------------------+
| Queries with Runtime in 95th Percentile |
+----------------------------------------+
1 row in set (0.05 sec)

...

mysql> CALL statement_performance_analyzer('delta', 'mydb.tmp_digests_ini', 'with_runtimes_in_95th_perc
+----------------------------------------+
| Next Output                            |
+----------------------------------------+
| Queries with Runtime in 95th Percentile |
+----------------------------------------+
1 row in set (0.03 sec)

...
```

Create an overall report of the 95th percentile queries and the top 10 queries with full table scans:

```
mysql> CALL statement_performance_analyzer('snapshot', NULL, NULL);
Query OK, 0 rows affected (0.01 sec)

mysql> SET @sys.statement_performance_analyzer.limit = 10;
Query OK, 0 rows affected (0.00 sec)

mysql> CALL statement_performance_analyzer('overall', NULL, 'with_runtimes_in_95th_percentile,with_full_tab
+---------------------------------------+
| Next Output                           |
+---------------------------------------+
| Queries with Runtime in 95th Percentile |
+---------------------------------------+
1 row in set (0.01 sec)

...

+-------------------------------------+
| Next Output                         |
+-------------------------------------+
| Top 10 Queries with Full Table Scan |
+-------------------------------------+
1 row in set (0.09 sec)

...
```

Use a custom view showing the top 10 queries sorted by total execution time, refreshing the view every minute using the `watch` command in Linux:

```
mysql> CREATE OR REPLACE VIEW mydb.my_statements AS
    -> SELECT sys.format_statement(DIGEST_TEXT) AS query,
    ->        SCHEMA_NAME AS db,
    ->        COUNT_STAR AS exec_count,
    ->        sys.format_time(SUM_TIMER_WAIT) AS total_latency,
    ->        sys.format_time(AVG_TIMER_WAIT) AS avg_latency,
    ->        ROUND(IFNULL(SUM_ROWS_SENT / NULLIF(COUNT_STAR, 0), 0)) AS rows_sent_avg,
    ->        ROUND(IFNULL(SUM_ROWS_EXAMINED / NULLIF(COUNT_STAR, 0), 0)) AS rows_examined_avg,
    ->        ROUND(IFNULL(SUM_ROWS_AFFECTED / NULLIF(COUNT_STAR, 0), 0)) AS rows_affected_avg,
    ->        DIGEST AS digest
    ->   FROM performance_schema.events_statements_summary_by_digest
    -> ORDER BY SUM_TIMER_WAIT DESC;
Query OK, 0 rows affected (0.01 sec)

mysql> CALL statement_performance_analyzer('create_table', 'mydb.digests_prev', NULL);
Query OK, 0 rows affected (0.10 sec)

shell> watch -n 60 "mysql sys --table -e \"
> SET @sys.statement_performance_analyzer.view = 'mydb.my_statements';
> SET @sys.statement_performance_analyzer.limit = 10;
> CALL statement_performance_analyzer('snapshot', NULL, NULL);
> CALL statement_performance_analyzer('delta', 'mydb.digests_prev', 'custom');
> CALL statement_performance_analyzer('save', 'mydb.digests_prev', NULL);
> \""

Every 60.0s: mysql sys --table -e "        ...  Mon Dec 22 10:58:51 2014

+---------------------------------+
| Next Output                     |
+---------------------------------+
| Top 10 Queries Using Custom View |
+---------------------------------+
+------------------+-------+-----------+---------------+-------------+---------------+-------------------
| query            | db    | exec_count | total_latency | avg_latency | rows_sent_avg | rows_examined_avg
+------------------+-------+-----------+---------------+-------------+---------------+-------------------
...
```

26.4.4.26 The table_exists() Procedure

Tests whether a given table exists as a regular table, a `TEMPORARY` table, or a view. The procedure returns the table type in an `OUT` parameter. If both a temporary and a permanent table exist with the given name, `TEMPORARY` is returned.

Parameters

- `in_db VARCHAR(64)`: The name of the database in which to check for table existance.
- `in_table VARCHAR(64)`: The name of the table to check the existance of.
- `out_exists ENUM('', 'BASE TABLE', 'VIEW', 'TEMPORARY')`: The return value. This is an `OUT` parameter, so it must be a variable into which the table type can be stored. When the procedure returns, the variable has one of the following values to indicate whether the table exists:
  - `''`: The table name does not exist as a base table, `TEMPORARY` table, or view.
  - `BASE TABLE`: The table name exists as a base (permanent) table.
  - `VIEW`: The table name exists as a view.
  - `TEMPORARY`: The table name exists as a `TEMPORARY` table.

Example

```
mysql> CREATE DATABASE db1;
Query OK, 1 row affected (0.01 sec)

mysql> USE db1;
Database changed
mysql> CREATE TABLE t1 (id INT PRIMARY KEY);
Query OK, 0 rows affected (0.03 sec)

mysql> CREATE TABLE t2 (id INT PRIMARY KEY);
Query OK, 0 rows affected (0.20 sec)

mysql> CREATE view v_t1 AS SELECT * FROM t1;
Query OK, 0 rows affected (0.02 sec)

mysql> CREATE TEMPORARY TABLE t1 (id INT PRIMARY KEY);
Query OK, 0 rows affected (0.00 sec)

mysql> CALL sys.table_exists('db1', 't1', @exists); SELECT @exists;
Query OK, 0 rows affected (0.01 sec)

+-----------+
| @exists   |
+-----------+
| TEMPORARY |
+-----------+
1 row in set (0.00 sec)

mysql> CALL sys.table_exists('db1', 't2', @exists); SELECT @exists;
Query OK, 0 rows affected (0.02 sec)

+------------+
| @exists    |
+------------+
| BASE TABLE |
+------------+
1 row in set (0.00 sec)

mysql> CALL sys.table_exists('db1', 'v_t1', @exists); SELECT @exists;
Query OK, 0 rows affected (0.02 sec)

+---------+
| @exists |
+---------+
| VIEW    |
+---------+
1 row in set (0.00 sec)

mysql> CALL sys.table_exists('db1', 't3', @exists); SELECT @exists;
Query OK, 0 rows affected (0.00 sec)
```

```
+---------+
| @exists |
+---------+
|         |
+---------+
1 row in set (0.00 sec)
```

26.4.5 sys Schema Stored Functions

The following sections describe `sys` schema stored functions.

26.4.5.1 The extract_schema_from_file_name() Function

Given a file path name, returns the path component that represents the schema name. This function assumes that the file name lies within the schema directory. For this reason, it will not work with partitions or tables defined using their own `DATA_DIRECTORY` table option.

This function is useful when extracting file I/O information from the Performance Schema that includes file path names. It provides a convenient way to display schema names, which can be more easily understood than full path names, and can be used in joins against object schema names.

Parameters

- `path VARCHAR(512)`: The full path to a data file from which to extract the schema name.

Return Value

A `VARCHAR(64)` value.

Example

```
mysql> SELECT extract_schema_from_file_name('/usr/local/mysql/data/world/City.ibd');
+-----------------------------------------------------------------------+
| extract_schema_from_file_name('/usr/local/mysql/data/world/City.ibd') |
+-----------------------------------------------------------------------+
| world                                                                 |
+-----------------------------------------------------------------------+
```

26.4.5.2 The extract_table_from_file_name() Function

Given a file path name, returns the path component that represents the table name.

This function is useful when extracting file I/O information from the Performance Schema that includes file path names. It provides a convenient way to display table names, which can be more easily understood than full path names, and can be used in joins against object table names.

Parameters

- `path VARCHAR(512)`: The full path to a data file from which to extract the table name.

Return Value

A `VARCHAR(64)` value.

Example

```
mysql> SELECT extract_table_from_file_name('/usr/local/mysql/data/world/City.ibd');
+-----------------------------------------------------------------------+
| extract_table_from_file_name('/usr/local/mysql/data/world/City.ibd') |
+-----------------------------------------------------------------------+
| City                                                                  |
```

```
+------------------------------------------------------------------------+
```

26.4.5.3 The format_bytes() Function

Given a value in bytes, converts it to human-readable format and returns a string consisting of a value and a units indicator. Depending on the size of the value, the units part is `bytes`, `KiB` (kibibytes), `MiB` (mebibytes), `GiB` (gibibytes), `TiB` (tebibytes), or `PiB` (pebibytes).

Parameters

- `bytes TEXT`: The bytes value to format.

Return Value

A `TEXT` value.

Example

```
mysql> SELECT format_bytes(512), format_bytes(18446644073709551615);
+------------------+-----------------------------------+
| format_bytes(512) | format_bytes(18446644073709551615) |
+------------------+-----------------------------------+
| 512 bytes         | 16383.91 PiB                      |
+------------------+-----------------------------------+
```

26.4.5.4 The format_path() Function

Given a path name, returns the modified path name after replacing subpaths that match the values of the following system variables, in order:

```
datadir
tmpdir
slave_load_tmpdir
innodb_data_home_dir
innodb_log_group_home_dir
innodb_undo_directory
basedir
```

A value that matches the value of system variable *`sysvar`* is replaced with the string `@@global.`*`sysvar`*.

Parameters

- `path VARCHAR(512)`: The path name to format.

Return Value

A `VARCHAR(512) CHARACTER SET utf8` value.

Example

```
mysql> SELECT format_path('/usr/local/mysql/data/world/City.ibd');
+-----------------------------------------------------+
| format_path('/usr/local/mysql/data/world/City.ibd') |
+-----------------------------------------------------+
| @@datadir/world/City.ibd                            |
+-----------------------------------------------------+
```

26.4.5.5 The format_statement() Function

Given a string (normally representing an SQL statement), reduces it to the length given by the `statement_truncate_len` configuration option, and returns the result. No truncation occurs if the

string is shorter than `statement_truncate_len`. Otherwise, the middle part of the string is replaced by an ellipsis (`...`).

This function is useful for formatting possibly lengthy statements retrieved from Performance Schema tables to a known fixed maximum length.

Parameters

- `statement LONGTEXT`: The statement to format.

Configuration Options

`format_statement()` operation can be modified using the following configuration options or their corresponding user-defined variables (see Section 26.4.2.1, “The sys_config Table”):

- `statement_truncate_len`, `@sys.statement_truncate_len`

  The maximum length of statements returned by the `format_statement()` function. Longer statements are truncated to this length. The default is 64.

Return Value

A `LONGTEXT` value.

Example

By default, `format_statement()` truncates statements to be no more than 64 characters. Setting `@sys.statement_truncate_len` changes the truncation length for the current session:

```
mysql> SET @stmt = 'SELECT variable, value, set_time, set_by FROM sys_config';
mysql> SELECT format_statement(@stmt);
+-----------------------------------------------------------+
| format_statement(@stmt)                                   |
+-----------------------------------------------------------+
| SELECT variable, value, set_time, set_by FROM sys_config |
+-----------------------------------------------------------+
mysql> SET @sys.statement_truncate_len = 32;
mysql> SELECT format_statement(@stmt);
+----------------------------------+
| format_statement(@stmt)          |
+----------------------------------+
| SELECT variabl ... ROM sys_config |
+----------------------------------+
```

26.4.5.6 The format_time() Function

Given a Performance Schema latency or wait time in picoseconds, converts it to human-readable format and returns a string consisting of a value and a units indicator. Depending on the size of the value, the units part is `ns` (nanoseconds), `us` (microseconds), `ms` (milliseconds), `s` (seconds), `m` (minutes), `h` (hours), `d` (days), or `w` (weeks).

Parameters

- `picoseconds TEXT`: The picoseconds value to format.

Return Value

A `TEXT` value.

Example

```
mysql> SELECT format_time(3501), format_time(188732396662000);
```

```
+------------------+-----------------------------+
| format_time(3501) | format_time(188732396662000) |
+------------------+-----------------------------+
| 3.50 ns          | 3.15 m                      |
+------------------+-----------------------------+
```

26.4.5.7 The list_add() Function

Adds a value to a comma-separated list of values and returns the result.

This function and `list_drop()` can be useful for manipulating the value of system variables such as `sql_mode` and `optimizer_switch` that take a comma-separated list of values.

Parameters

- `in_list TEXT`: The list to be modified.
- `in_add_value TEXT`: The value to add to the list.

Return Value

A `TEXT` value.

Example

```
mysql> SELECT @@sql_mode;
+----------------------------------------+
| @@sql_mode                             |
+----------------------------------------+
| ONLY_FULL_GROUP_BY,STRICT_TRANS_TABLES |
+----------------------------------------+
mysql> SET @@sql_mode = list_add(@@sql_mode, 'NO_ENGINE_SUBSTITUTION');
mysql> SELECT @@sql_mode;
+-----------------------------------------------------------------+
| @@sql_mode                                                      |
+-----------------------------------------------------------------+
| ONLY_FULL_GROUP_BY,STRICT_TRANS_TABLES,NO_ENGINE_SUBSTITUTION |
+-----------------------------------------------------------------+
mysql> SET @@sql_mode = list_drop(@@sql_mode, 'ONLY_FULL_GROUP_BY');
mysql> SELECT @@sql_mode;
+-----------------------------------------+
| @@sql_mode                              |
+-----------------------------------------+
| STRICT_TRANS_TABLES,NO_ENGINE_SUBSTITUTION |
+-----------------------------------------+
```

26.4.5.8 The list_drop() Function

Removes a value from a comma-separated list of values and returns the result. For more information, see the description of `list_add()`

Parameters

- `in_list TEXT`: The list to be modified.
- `in_drop_value TEXT`: The value to drop from the list.

Return Value

A `TEXT` value.

26.4.5.9 The ps_is_account_enabled() Function

Returns `YES` or `NO` to indicate whether Performance Schema instrumentation for a given account is enabled.

Parameters

- `in_host VARCHAR(60)`: The host name of the account to check.
- `in_user VARCHAR(32)`: The user name of the account to check.

Return Value

An `ENUM('YES','NO')` value.

Example

```
mysql> SELECT ps_is_account_enabled('localhost', 'root');
+-----------------------------------------+
| ps_is_account_enabled('localhost', 'root') |
+-----------------------------------------+
| YES                                     |
+-----------------------------------------+
```

26.4.5.10 The ps_is_consumer_enabled() Function

Returns `YES` or `NO` to indicate whether a given Performance Schema consumer is enabled, or `NULL` if the argument is not a valid consumer name.

This function accounts for the consumer hierarchy, so a consumer is not considered enabled unless all consumers on which depends are also enabled. For information about the consumer hierarchy, see Section 25.4.7, “Pre-Filtering by Consumer”.

Parameters

- `in_consumer VARCHAR(64)`: The name of the consumer to check.

Return Value

An `ENUM('YES','NO')` value.

Example

```
mysql> SELECT ps_is_consumer_enabled('thread_instrumentation');
+-------------------------------------------------+
| ps_is_consumer_enabled('thread_instrumentation') |
+-------------------------------------------------+
| YES                                             |
+-------------------------------------------------+
```

26.4.5.11 The ps_is_instrument_default_enabled() Function

Returns `YES` or `NO` to indicate whether a given Performance Schema instrument is enabled by default.

Parameters

- `in_instrument VARCHAR(128)`: The name of the instrument to check.

Return Value

An `ENUM('YES','NO')` value.

Example

```
mysql> SELECT ps_is_instrument_default_enabled('memory/innodb/row_log_buf');
+-------------------------------------------------------------+
```

```
| ps_is_instrument_default_enabled('memory/innodb/row_log_buf') |
+---------------------------------------------------------------+
| NO                                                            |
+---------------------------------------------------------------+
mysql> SELECT ps_is_instrument_default_enabled('statement/sql/alter_user');
+---------------------------------------------------------------+
| ps_is_instrument_default_enabled('statement/sql/alter_user') |
+---------------------------------------------------------------+
| YES                                                           |
+---------------------------------------------------------------+
```

26.4.5.12 The ps_is_instrument_default_timed() Function

Returns `YES` or `NO` to indicate whether a given Performance Schema instrument is timed by default.

Parameters

- `in_instrument VARCHAR(128)`: The name of the instrument to check.

Return Value

An `ENUM('YES','NO')` value.

Example

```
mysql> SELECT ps_is_instrument_default_timed('memory/innodb/row_log_buf');
+---------------------------------------------------------------+
| ps_is_instrument_default_timed('memory/innodb/row_log_buf') |
+---------------------------------------------------------------+
| NO                                                            |
+---------------------------------------------------------------+
mysql> SELECT ps_is_instrument_default_timed('statement/sql/alter_user');
+---------------------------------------------------------------+
| ps_is_instrument_default_timed('statement/sql/alter_user') |
+---------------------------------------------------------------+
| YES                                                           |
+---------------------------------------------------------------+
```

26.4.5.13 The ps_is_thread_instrumented() Function

Returns `YES` or `NO` to indicate whether Performance Schema instrumentation for a given connection ID is enabled, `UNKNOWN` if the ID is unknown, or `NULL` if the ID is `NULL`.

Parameters

- `in_connection_id BIGINT UNSIGNED`: The connection ID. This is a connection ID as given in the `PROCESSLIST_ID` column of the Performance Schema `threads` table or the `Id` column of `SHOW PROCESSLIST` output.

Return Value

An `ENUM('YES','NO','UNKNOWN')` value.

Example

```
mysql> SELECT ps_is_thread_instrumented(43);
+-------------------------------+
| ps_is_thread_instrumented(43) |
+-------------------------------+
| UNKNOWN                       |
+-------------------------------+
mysql> SELECT ps_is_thread_instrumented(CONNECTION_ID());
+-----------------------------------------+
| ps_is_thread_instrumented(CONNECTION_ID()) |
```

```
+-------------------------------------------+
| YES                                       |
+-------------------------------------------+
```

26.4.5.14 The ps_thread_account() Function

Given a Performance Schema thread ID, returns the *`user_name@host_name`* account associated with the thread.

Parameters

- `in_thread_id BIGINT UNSIGNED`: The thread ID for which to return the account. The value should match the `THREAD_ID` column from some Performance Schema `threads` table row.

Return Value

A `TEXT` value.

Example

```
mysql> SELECT ps_thread_account(ps_thread_id(CONNECTION_ID()));
+-------------------------------------------------+
| ps_thread_account(ps_thread_id(CONNECTION_ID())) |
+-------------------------------------------------+
| root@localhost                                  |
+-------------------------------------------------+
```

26.4.5.15 The ps_thread_id() Function

Returns the Performance Schema thread ID for a given connection ID, or the thread ID for the current connection if the connection ID is `NULL`.

Parameters

- `in_connection_id BIGINT UNSIGNED`: The ID of the connection for which to return the thread ID. This is a connection ID as given in the `PROCESSLIST_ID` column of the Performance Schema `threads` table or the `Id` column of `SHOW PROCESSLIST` output.

Return Value

A `BIGINT UNSIGNED` value.

Example

```
mysql> SELECT ps_thread_id(260);
+-------------------+
| ps_thread_id(260) |
+-------------------+
|               285 |
+-------------------+
```

26.4.5.16 The ps_thread_stack() Function

Returns a JSON formatted stack of all statements, stages, and events within the Performance Schema for a given thread ID.

Parameters

- `in_thread_id BIGINT`: The ID of the thread to trace. The value should match the `THREAD_ID` column from some Performance Schema `threads` table row.
- `in_verbose BOOLEAN`: Whether to include `file:lineno` information in the events.

Return Value

A `LONGTEXT CHARACTER SET latin1` value.

Example

```
mysql> SELECT ps_thread_stack(37, FALSE) AS thread_stack\G
*************************** 1. row ***************************
thread_stack: {"rankdir": "LR","nodesep": "0.10",
"stack_created": "2014-02-19 13:39:03", "mysql_version": "8.0.2-dmr-debug-log",
"mysql_user": "root@localhost","events": [{"nesting_event_id": "0",
"event_id": "10", "timer_wait": 256.35, "event_info": "sql/select",
"wait_info": "select @@version_comment limit 1\nerrors: 0\nwarnings: 0\nlock time:
...
```

26.4.5.17 The ps_thread_trx_info() Function

Returns a JSON object containing information about a given thread. The information includes the current transaction, and the statements it has already executed, derived from the Performance Schema `events_transactions_current` and `events_statements_history` tables. (The consumers for those tables must be enabled to obtain full data in the JSON object.)

If the output exceeds the truncation length (65535 by default), a JSON error object is returned, such as:

```
{ "error": "Trx info truncated: Row 6 was cut by GROUP_CONCAT()" }
```

Similar error objects are returned for other warnings and exceptions raised during function execution.

Parameters

- `in_thread_id BIGINT UNSIGNED`: The thread ID for which to return transaction information. The value should match the `THREAD_ID` column from some Performance Schema `threads` table row.

Configuration Options

`ps_thread_trx_info()` operation can be modified using the following configuration options or their corresponding user-defined variables (see Section 26.4.2.1, “The sys_config Table”):

- `ps_thread_trx_info.max_length`, `@sys.ps_thread_trx_info.max_length`

  The maximum length of the output. The default is 65535.

Return Value

A `LONGTEXT` value.

Example

```
mysql> SELECT ps_thread_trx_info(48)\G
*************************** 1. row ***************************
ps_thread_trx_info(48): [
  {
    "time": "790.70 us",
    "state": "COMMITTED",
    "mode": "READ WRITE",
    "autocommitted": "NO",
    "gtid": "AUTOMATIC",
    "isolation": "REPEATABLE READ",
    "statements_executed": [
      {
        "sql_text": "INSERT INTO info VALUES (1, \'foo\')",
        "time": "471.02 us",
        "schema": "trx",
        "rows_examined": 0,
```

```
        "rows_affected": 1,
        "rows_sent": 0,
        "tmp_tables": 0,
        "tmp_disk_tables": 0,
        "sort_rows": 0,
        "sort_merge_passes": 0
      },
      {
        "sql_text": "COMMIT",
        "time": "254.42 us",
        "schema": "trx",
        "rows_examined": 0,
        "rows_affected": 0,
        "rows_sent": 0,
        "tmp_tables": 0,
        "tmp_disk_tables": 0,
        "sort_rows": 0,
        "sort_merge_passes": 0
      }
    ]
  },
  {
    "time": "426.20 us",
    "state": "COMMITTED",
    "mode": "READ WRITE",
    "autocommitted": "NO",
    "gtid": "AUTOMATIC",
    "isolation": "REPEATABLE READ",
    "statements_executed": [
      {
        "sql_text": "INSERT INTO info VALUES (2, \'bar\')",
        "time": "107.33 us",
        "schema": "trx",
        "rows_examined": 0,
        "rows_affected": 1,
        "rows_sent": 0,
        "tmp_tables": 0,
        "tmp_disk_tables": 0,
        "sort_rows": 0,
        "sort_merge_passes": 0
      },
      {
        "sql_text": "COMMIT",
        "time": "213.23 us",
        "schema": "trx",
        "rows_examined": 0,
        "rows_affected": 0,
        "rows_sent": 0,
        "tmp_tables": 0,
        "tmp_disk_tables": 0,
        "sort_rows": 0,
        "sort_merge_passes": 0
      }
    ]
  }
]
```


26.4.5.18 The quote_identifier() Function

Given a string argument, this function produces a quoted identifier suitable for inclusion in SQL statements. This is useful when a value to be used as an identifier is a reserved word or contains backtick (`) characters.

Parameters

`in_identifier TEXT`: The identifier to quote.

Return Value

A `TEXT` value.

Example

```
mysql> SELECT quote_identifier('plain');
+--------------------------+
| quote_identifier('plain') |
+--------------------------+
| `plain`                  |
+--------------------------+
mysql> SELECT quote_identifier('trick`ier');
+-----------------------------+
| quote_identifier('trick`ier') |
+-----------------------------+
| `trick``ier`                |
+-----------------------------+
mysql> SELECT quote_identifier('integer');
+----------------------------+
| quote_identifier('integer') |
+----------------------------+
| `integer`                  |
+----------------------------+
```

26.4.5.19 The sys_get_config() Function

Given a configuration option name, returns the option value from the `sys_config` table, or the provided default value (which may be `NULL`) if the option does not exist in the table.

If `sys_get_config()` returns the default value and that value is `NULL`, it is expected that the caller is able to handle `NULL` for the given configuration option.

By convention, routines that call `sys_get_config()` first check whether the corresponding user-defined variable exists and is non-`NULL`. If so, the routine uses the variable value without reading the `sys_config` table. If the variable does not exist or is `NULL`, the routine reads the option value from the table and sets the user-defined variable to that value. For more information about the relationship between configuration options and their corresponding user-defined variables, see Section 26.4.2.1, “The sys_config Table”.

If you want to check whether the configuration option has already been set and, if not, use the return value of `sys_get_config()`, you can use `IFNULL(...)` (see example later). However, this should not be done inside a loop (for example, for each row in a result set) because for repeated calls where the assignment is needed only in the first iteration, using `IFNULL(...)` is expected to be significantly slower than using an `IF (...) THEN ... END IF;` block (see example later).

Parameters

- `in_variable_name VARCHAR(128)`: The name of the configuration option for which to return the value.
- `in_default_value VARCHAR(128)`: The default value to return if the configuration option is not found in the `sys_config` table.

Return Value

A `VARCHAR(128)` value.

Example

Get a configuration value from the `sys_config` table, falling back to 128 as the default if the option is not present in the table:

```
mysql> SELECT sys.sys_get_config('statement_truncate_len', 128) AS Value;
+-------+
| Value |
+-------+
```

```
| 64    |
+-------+
```

One-liner example: Check whether the option is already set; if not, assign the `IFNULL(...)` result (using the value from the `sys_config` table):

```
mysql> SET @sys.statement_truncate_len =
    -> IFNULL(@sys.statement_truncate_len,
    ->        sys.sys_get_config('statement_truncate_len', 64));
```

`IF (...) THEN ... END IF;` block example: Check whether the option is already set; if not, assign the value from the `sys_config` table:

```
IF (@sys.statement_truncate_len IS NULL) THEN
  SET @sys.statement_truncate_len = sys.sys_get_config('statement_truncate_len', 64);
END IF;
```

26.4.5.20 The version_major() Function

This function returns the major version of the MySQL server.

Parameters

None.

Return Value

A `TINYINT UNSIGNED` value.

Example

```
mysql> SELECT VERSION(), version_major();
+---------------------+-----------------+
| VERSION()           | version_major() |
+---------------------+-----------------+
| 8.0.2-dmr-debug-log |               8 |
+---------------------+-----------------+
```

26.4.5.21 The version_minor() Function

This function returns the minor version of the MySQL server.

Parameters

None.

Return Value

A `TINYINT UNSIGNED` value.

Example

```
mysql> SELECT VERSION(), version_minor();
+---------------------+-----------------+
| VERSION()           | version_minor() |
+---------------------+-----------------+
| 8.0.2-dmr-debug-log |               0 |
+---------------------+-----------------+
```

26.4.5.22 The version_patch() Function

This function returns the patch release version of the MySQL server.

Parameters

None.

Return Value

A `TINYINT UNSIGNED` value.

Example

```
mysql> SELECT VERSION(), version_patch();
+---------------------+-----------------+
| VERSION()           | version_patch() |
+---------------------+-----------------+
| 8.0.2-dmr-debug-log |               2 |
+---------------------+-----------------+
```

Chapter 27 Connectors and APIs

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MySQL Connectors provide connectivity to the MySQL server for client programs. APIs provide low-level access to the MySQL protocol and MySQL resources. Both Connectors and the APIs enable you to connect and execute MySQL statements from another language or environment, including ODBC, Java (JDBC), Perl, Python, PHP, Ruby, and native C MySQL instances.

Note

Connector version numbers do not correlate with MySQL Server version numbers. See Table 27.2, “MySQL Connector Versions and MySQL Server Versions”.

MySQL Connectors

Oracle develops a number of connectors:

- Connector/C is a standalone replacement for the MySQL Client Library (`libmysqlclient`), to be used for C applications.
- Connector/C++ enables C++ applications to connect to MySQL.
- Connector/J provides driver support for connecting to MySQL from Java applications using the standard Java Database Connectivity (JDBC) API.
- Connector/Net enables developers to create .NET applications that connect to MySQL. Connector/Net implements a fully functional ADO.NET interface and provides support for use with ADO.NET aware tools. Applications that use Connector/Net can be written in any supported .NET language.

  MySQL for Visual Studio works with Connector/Net and Microsoft Visual Studio 2012, 2013, 2015, and 2017. MySQL for Visual Studio provides access to MySQL objects and data from Visual Studio. As a Visual Studio package, it integrates directly into Server Explorer providing the ability to create new connections and work with MySQL database objects.
- Connector/ODBC provides driver support for connecting to MySQL using the Open Database Connectivity (ODBC) API. Support is available for ODBC connectivity from Windows, Unix, and OS X platforms.
- Connector/Python provides driver support for connecting to MySQL from Python applications using an API that is compliant with the Python DB API version 2.0. No additional Python modules or MySQL client libraries are required.

The MySQL C API

For direct access to using MySQL natively within a C application, the C API provides low-level access to the MySQL client/server protocol through the `libmysqlclient` client library. This is the primary method used to connect to an instance of the MySQL server, and is used both by MySQL command-line clients and many of the MySQL Connectors and third-party APIs detailed here.

`libmysqlclient` is included in MySQL distributions and in Connector/C distributions.

See also Section 27.7.1, “MySQL C API Implementations”.

To access MySQL from a C application, or to build an interface to MySQL for a language not supported by the Connectors or APIs in this chapter, the C API is where to start. A number of programmer's utilities are available to help with the process; see Section 4.7, “MySQL Program Development Utilities”.

Third-Party MySQL APIs

The remaining APIs described in this chapter provide an interface to MySQL from specific application languages. These third-party solutions are not developed or supported by Oracle. Basic information on their usage and abilities is provided here for reference purposes only.

All the third-party language APIs are developed using one of two methods, using `libmysqlclient` or by implementing a *native driver*. The two solutions offer different benefits:

- Using *`libmysqlclient`* offers complete compatibility with MySQL because it uses the same libraries as the MySQL client applications. However, the feature set is limited to the implementation and interfaces exposed through `libmysqlclient` and the performance may be lower as data is copied between the native language, and the MySQL API components.
- *Native drivers* are an implementation of the MySQL network protocol entirely within the host language or environment. Native drivers are fast, as there is less copying of data between components, and they can offer advanced functionality not available through the standard MySQL API. Native drivers are also easier for end users to build and deploy because no copy of the MySQL client libraries is needed to build the native driver components.

Table 27.1, “MySQL APIs and Interfaces” lists many of the libraries and interfaces available for MySQL. Table 27.2, “MySQL Connector Versions and MySQL Server Versions” shows which MySQL Server versions each connector supports.

Table 27.1 MySQL APIs and Interfaces

Environment	API	Type	Notes
Ada	GNU Ada MySQL Bindings	`libmysqlclient`	See MySQL Bindings for GNU Ada
C	C API	`libmysqlclient`	See Section 27.7, “MySQL C API”.
C	Connector/C	Replacement for `libmysqlclient`	See MySQL Connector/C Developer Guide.
C++	Connector/C++	`libmysqlclient`	See MySQL Connector/C++ Developer Guide.
	MySQL++	`libmysqlclient`	See MySQL++ Web site.
	MySQL wrapped	`libmysqlclient`	See MySQL wrapped.
Cocoa	MySQL-Cocoa	`libmysqlclient`	Compatible with the Objective-C Cocoa environment. See http://mysql-cocoa.sourceforge.net/
D	MySQL for D	`libmysqlclient`	See MySQL for D.
Eiffel	Eiffel MySQL	`libmysqlclient`	See Section 27.13, “MySQL Eiffel Wrapper”.
Erlang	`erlang-mysql-driver`	`libmysqlclient`	See `erlang-mysql-driver`.
Haskell	Haskell MySQL Bindings	Native Driver	See Brian O'Sullivan's pure Haskell MySQL bindings.
	`hsql-mysql`	`libmysqlclient`	See MySQL driver for Haskell .
Java/ JDBC	Connector/J	Native Driver	See MySQL Connector/J 5.1 Developer Guide.
Kaya	MyDB	`libmysqlclient`	See MyDB.
Lua	LuaSQL	`libmysqlclient`	See LuaSQL.
.NET/ Mono	Connector/Net	Native Driver	See MySQL Connector/Net Developer Guide.
Objective Caml	OBjective Caml MySQL Bindings	`libmysqlclient`	See MySQL Bindings for Objective Caml.
Octave	Database bindings for GNU Octave	`libmysqlclient`	See Database bindings for GNU Octave.
ODBC	Connector/ODBC	`libmysqlclient`	See MySQL Connector/ODBC Developer Guide.
Perl	`DBI`/`DBD::mysql`	`libmysqlclient`	See Section 27.9, “MySQL Perl API”.
	`Net::MySQL`	Native Driver	See `Net::MySQL` at CPAN
PHP	`mysql`, `ext/mysql` interface (deprecated)	`libmysqlclient`	See Original MySQL API.
	`mysqli`, `ext/mysqli` interface	`libmysqlclient`	See MySQL Improved Extension.
	`PDO_MYSQL`	`libmysqlclient`	See MySQL Functions (PDO_MYSQL).
	PDO mysqlnd	Native Driver	
Python	Connector/Python	Native Driver	See MySQL Connector/Python Developer Guide.

Environment	API	Type	Notes
Python	Connector/Python C Extension	`libmysqlclient`	See MySQL Connector/Python Developer Guide.
	MySQLdb	`libmysqlclient`	See Section 27.10, “MySQL Python API”.
Ruby	MySQL/Ruby	`libmysqlclient`	Uses `libmysqlclient`. See Section 27.11.1, “The MySQL/Ruby API”.
	Ruby/MySQL	Native Driver	See Section 27.11.2, “The Ruby/MySQL API”.
Scheme	`Myscsh`	`libmysqlclient`	See `Myscsh`.
SPL	`sql_mysql`	`libmysqlclient`	See `sql_mysql` for SPL.
Tcl	MySQLtcl	`libmysqlclient`	See Section 27.12, “MySQL Tcl API”.

Table 27.2 MySQL Connector Versions and MySQL Server Versions

Connector	Connector version	MySQL Server version
Connector/C	6.1.0 GA	5.6, 5.5, 5.1, 5.0, 4.1
Connector/C++	1.0.5 GA	5.6, 5.5, 5.1
Connector/J	5.1.8	5.6, 5.5, 5.1, 5.0, 4.1
Connector/Net	6.9.9 GA	5.7, 5.6, 5.5, 5.1, 5.0
Connector/Net	6.8.8 GA	5.7, 5.6, 5.5, 5.1, 5.0
Connector/ODBC	5.1	5.6, 5.5, 5.1, 5.0, 4.1.1+
Connector/ODBC	3.51 (Unicode not supported)	5.6, 5.5, 5.1, 5.0, 4.1
Connector/Python	2.0	5.7, 5.6, 5.5
Connector/Python	1.2	5.7, 5.6, 5.5

27.1 MySQL Connector/C

The MySQL Connector/C manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/C Developer Guide
- Release notes: MySQL Connector/C Release Notes

27.2 MySQL Connector/C++

The MySQL Connector/C++ manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/C++ Developer Guide
- Release notes: MySQL Connector/C++ Release Notes

27.3 MySQL Connector/J

The MySQL Connector/J manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/J 5.1 Developer Guide
- Release notes: MySQL Connector/J Release Notes

27.4 MySQL Connector/Net

The MySQL Connector/Net manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/Net Developer Guide
- Release notes: MySQL Connector/Net Release Notes

27.5 MySQL Connector/ODBC

The MySQL Connector/ODBC manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/ODBC Developer Guide
- Release notes: MySQL Connector/ODBC Release Notes

27.6 MySQL Connector/Python

The MySQL Connector/Python manual is published in standalone form, not as part of the MySQL Reference Manual. For information, see these documents:

- Main manual: MySQL Connector/Python Developer Guide
- Release notes: MySQL Connector/Python Release Notes

27.7 MySQL C API

The C API provides low-level access to the MySQL client/server protocol and enables C programs to access database contents. The C API code is distributed with MySQL and implemented in the `libmysqlclient` library. See Section 27.7.1, “MySQL C API Implementations”.

Most other client APIs use the `libmysqlclient` library to communicate with the MySQL server. (Exceptions are Connector/J and Connector/Net.) This means that, for example, you can take advantage of many of the same environment variables that are used by other client programs because they are referenced from the library. For a list of these variables, see Section 4.1, “Overview of MySQL Programs”.

For instructions on building client programs using the C API, see Section 27.7.4.1, “Building C API Client Programs”. For programming with threads, see Section 27.7.4.3, “Writing C API Threaded Client Programs”.

Note

If, after an upgrade, you experience problems with compiled client programs, such as `Commands out of sync` or unexpected core dumps, the programs were probably compiled using old header or library files. In this case, check the date of the `mysql.h` file and `libmysqlclient.a` library used for compilation to verify that they are from the new MySQL distribution. If not, recompile the programs with the new headers and libraries. Recompilation might also be necessary for programs compiled against the shared client library if the library major version number has changed (for example, from `libmysqlclient.so.17` to `libmysqlclient.so.18`). For additional compatibility information, see Section 27.7.4.4, “Running C API Client Programs”.

Clients have a maximum communication buffer size. The size of the buffer that is allocated initially (16KB) is automatically increased up to the maximum size (16MB by default). Because buffer sizes are increased only as demand warrants, simply increasing the maximum limit does not in itself cause

more resources to be used. This size check is mostly a precaution against erroneous statements and communication packets.

The communication buffer must be large enough to contain a single SQL statement (for client-to-server traffic) and one row of returned data (for server-to-client traffic). Each session's communication buffer is dynamically enlarged to handle any query or row up to the maximum limit. For example, if you have `BLOB` values that contain up to 16MB of data, you must have a communication buffer limit of at least 16MB (in both server and client). The default maximum built into the client library is 1GB, but the default maximum in the server is 1MB. You can increase this by changing the value of the `max_allowed_packet` parameter at server startup. See Section 5.1.1, “Configuring the Server”.

The MySQL server shrinks each communication buffer to `net_buffer_length` bytes after each query. For clients, the size of the buffer associated with a connection is not decreased until the connection is closed, at which time client memory is reclaimed.

27.7.1 MySQL C API Implementations

The MySQL C API is a C-based API that client applications written in C can use to communicate with MySQL Server. Client programs refer to C API header files at compile time and link to a C API library file, `libmysqlclient`, at link time.

There are two ways to obtain the C API header and library files required to build C API client programs:

- Install a MySQL Server distribution.
- Install a Connector/C distribution.

For both MySQL Server and Connector/C, you can install a binary distribution that contains the C API files pre-built, or you can use a source distribution and build the C API files yourself.

Normally, you install either a MySQL Server distribution or a Connector/C distribution, but not both. For information about issues involved with simultaneous MySQL Server and Connector/C installations, see Section 27.7.2, “Simultaneous MySQL Server and Connector/C Installations”.

The names of the library files to use when linking C API client applications depend on the library type and platform for which a distribution is built:

- On Unix (and Unix-like) sytems, the static library is `libmysqlclient.a`. The dynamic library is `libmysqlclient.so` on most Unix systems and `libmysqlclient.dylib` on OS X.
- On Windows, the static library is `mysqlclient.lib` and the dynamic library is `libmysql.dll`. Windows distributions also include `libmysql.lib`, a static import library needed for using the dynamic library.

  Windows distributions also include a set of debug libraries. These have the same names as the nondebug libraries, but are located in the `lib/debug` library. You must use the debug libraries when compiling clients built using the debug C runtime.

On Unix, you may also see libraries that include `_r` in the names. Before MySQL 5.5, these were built as thread-safe (re-entrant) libraries separately from the non-`_r` libraries. As of 5.5, both libraries are the same and the `_r` names are symbolic links to the corresponding non-`_r` names. There is no need to use the `_r` libraries. For example, if you use `mysql_config` to obtain linker flags, you can use `mysql_config --libs` in all cases, even for threaded clients. There is no need to use `mysql_config --libs_r`.

27.7.2 Simultaneous MySQL Server and Connector/C Installations

MySQL Server and Connector/C installation packages both provide the files needed to build and run MySQL C API client programs. This section discusses when it is possible to install both products on the same system. For some packaging formats, this is possible without conflict. For others, both products cannot be installed at the same time.

This discussion assumes the use of similar package types for both products (for example, RPM packages for both products). It does not try to describe coexistence between packaging types (for example, use of RPM packages for one product and a `tar` file package for the other). Nor does it describe coexistence of packages provided by Oracle and those provided by third-party vendors.

If you install both products, it may be necessary to adjust your development tools or runtime environment to choose one set of header files and libraries over the other. See Section 27.7.4.1, "Building C API Client Programs", and Section 27.7.4.4, "Running C API Client Programs".

`tar` and Zip file packages install under the directory into which you unpack them. For example, you can unpack MySQL Server and Connector/C `tar` packages under `/usr/local` and they will unpack into distinct directory names without conflict.

Windows MSI installers use their own installation directory, so MySQL Server and Connector/C installers do not conflict.

OS X DMG packages install under the same parent directory but in a different subdirectory, so there is no conflict. For example:

```
/usr/local/mysql-5.6.11-osx10.7-x86_64/
/usr/local/mysql-connector-c-6.1.0-osx10.7-x86/
```

Solaris PKG packages install under the same parent directory but in a different subdirectory, so there is no conflict. For example:

```
/opt/mysql/mysql
/opt/mysql/connector-c
```

The Solaris Connector/C installer does not create any symlinks from system directories such as `/usr/bin` or `/usr/lib` into the installation directory. That must be done manually if desired after installation.

For RPM installations, there are several types of RPM packages. MySQL Server `shared` and `devel` RPM packages are similar to the corresponding Connector/C RPM packages. These RPM package types cannot coexist because the MySQL Server and Connector/C RPM packages use the same installation locations for the client library-related files. This means the following conditions hold:

- If MySQL Server `shared` and `devel` RPM packages are installed, they provide the C API headers and libraries, and there is no need to install the Connector/C RPM packages. To install the Connector/C packages anyway, you must first remove the corresponding MySQL Server packages.
- To install MySQL Server RPM packages if you already have Connector/C RPM packages installed, you must first remove the Connector/C RPM packages.

MySQL Server RPM packages other than `shared` and `devel` do not conflict with Connector/C packages and can be installed if Connector/C is installed. This includes the main server RPM that includes the `mysqld` server itself.

27.7.3 Example C API Client Programs

Many of the clients in MySQL source distributions are written in C, such as `mysql`, `mysqladmin`, and `mysqlshow`. If you are looking for examples that demonstrate how to use the C API, take a look at these clients: Obtain a source distribution and look in its `client` directory. See Section 2.1.2, "How to Get MySQL".

27.7.4 Building and Running C API Client Programs

The following sections provide information on building client programs that use the C API. Topics include compiling and linking clients, writing threaded clients, and troubleshooting runtime problems.

27.7.4.1 Building C API Client Programs

This section provides guidelines for compiling C programs that use the MySQL C API.

Compiling MySQL Clients on Unix

The examples here use `gcc` as the compiler. A different compiler might be appropriate on some systems (for example, `clang` on OS X or FreeBSD, or Sun Studio on Solaris). Adjust the examples as necessary.

You may need to specify an `-I` option when you compile client programs that use MySQL header files, so that the compiler can find them. For example, if the header files are installed in `/usr/local/mysql/include`, use this option in the compile command:

```
-I/usr/local/mysql/include
```

MySQL clients must be linked using the `-lmysqlclient` option in the link command. You may also need to specify a `-L` option to tell the linker where to find the library. For example, if the library is installed in `/usr/local/mysql/lib`, use these options in the link command:

```
-L/usr/local/mysql/lib -lmysqlclient
```

The path names may differ on your system. Adjust the `-I` and `-L` options as necessary.

To make it simpler to compile MySQL programs on Unix, use the `mysql_config` script. See Section 4.7.1, “`mysql_config` — Display Options for Compiling Clients”.

`mysql_config` displays the options needed for compiling or linking:

```
shell> mysql_config --cflags
shell> mysql_config --libs
```

You can run those commands to get the proper options and add them manually to compilation or link commands. Alternatively, include the output from `mysql_config` directly within command lines using backticks:

```
shell> gcc -c `mysql_config --cflags` progname.c
shell> gcc -o progname progname.o `mysql_config --libs`
```

On Unix, linking uses dynamic libraries by default. To link to the static client library instead, add its path name to the link command. For example, if the library is located in `/usr/local/mysql/lib`, link like this:

```
shell> gcc -o progname progname.o /usr/local/mysql/lib/libmysqlclient.a
```

Or use `mysql_config` to provide the library name:

```
shell> gcc -o progname progname.o `mysql_config --variable=pkglibdir`/libmysqlclient.a
```

`mysql_config` does not currently provide a way to list all libraries needed for static linking, so it might be necessary to name additional libraries on the link command (for example, `-lnsl -lsocket` on Solaris). To get an idea which libraries to add, use `mysql_config --libs` and `ldd libmysqlclient.so` (or `otool -L libmysqlclient.dylib` on OS X).

`pkg-config` can be used as an alternative to `mysql_config` for obtaining information such as compiler flags or link libraries required to compile MySQL applications. For example, the following pairs of commands are equivalent:

```
mysql_config --cflags
pkg-config --cflags mysqlclient

mysql_config --libs
pkg-config --libs mysqlclient
```

To produce flags for static linking, use this command:

```
pkg-config --static --libs mysqlclient
```

For more information, see Section 27.7.4.2, “Building C API Client Programs Using pkg-config”.

Compiling MySQL Clients on Microsoft Windows

To specify header and library file locations, use the facilities provided by your development environment.

To build C API clients on Windows, you must link in the C client library, as well as the Windows ws2_32 sockets library and Secur32 security library.

You link your code with either the dynamic or static C client library. On Windows, the static library is named `mysqlclient.lib` and the dynamic library is named `libmysql.dll`. In addition, the `libmysql.lib` static import library is needed for using the dynamic library. If the static C client library is used, the client application must be compiled with the same version of Visual Studio used to compile the C client library (which is Visual Studio 2015 for the static C client library built by Oracle).

Note

The MySQL Connector/C is a standalone, drop-in replacement of the MySQL C client libraries that come with the MySQL server distribution. The Oracle-built MySQL Connector/C contains currently two versions of the static client library, one built with Visual Studio 2013 and the other one with Visual Studio 2015; use the one that matches the Visual Studio version you use to compile your application.

When using the Oracle-built MySQL C client library (or MySQL Connector/C), following these rules when it comes to linking the C runtime for your client application:

- For the Community version of the MySQL C client library (or the Community version of MySQL Connector/C):
  - *For version 8.0.0* (or MySQL Connector/C Community 6.1.9 and before):
    - If linking to the static C client library, link statically to the C runtime (use the `/MT` compiler option).
    - If linking to the dynamic C client library, link either statically or dynamically to the C runtime (use either `/MT` or `/MD` compiler option).
  - *For version 8.0.1 and later* (or MySQL Connector/C Community 6.1.10 and later): Always link dynamically to the C runtime (use the `/MD` compiler option), whether you are linking to the static or dynamic C client library. Also, target hosts running the client application need to have the Visual C++ Redistributable for Visual Studio 2015 installed.
- For the Commercial version of the MySQL C client library (or the Commercial version of MySQL Connector/C):
  - If linking to the static C client library, link statically to the C runtime (use the `/MT` compiler option).
  - If linking to the dynamic C client library, link either statically or dynamically to the C runtime (use either `/MT` or `/MD` compiler option).

In general, when linking to a static MySQL C client library, the client library and the client application must use the same compiler option when it comes to linking the C runtime—that is, if your C client library is compiled with the `/MT` option, your client application should also be compiled with the `/MT` option, and so on (see the MSDN page describing the C library linking options for more details). Follow this rule when you are building your own static MySQL C client library (or MySQL Connector/C) from source and linking you client application to it.

Note

Debug Mode: Because of the above-mentioned rule, you cannot build your application in debug mode (with the `/MTd` or `/MDd` compiler option) and link it to the static C client library built by Oracle, which is *not* built with the debug options; instead, you will have to build the static client library from source with the debug options.

Troubleshooting Problems Linking to the MySQL Client Library

The MySQL client library includes SSL support built in. It is unnecessary to specify either `-lssl` or `-lcrypto` at link time. Doing so may in fact result in problems at runtime.

If the linker cannot find the MySQL client library, you might get undefined-reference errors for symbols that start with `mysql_`, such as those shown here:

```
/tmp/ccFKsdPa.o: In function `main':
/tmp/ccFKsdPa.o(.text+0xb): undefined reference to `mysql_init'
/tmp/ccFKsdPa.o(.text+0x31): undefined reference to `mysql_real_connect'
/tmp/ccFKsdPa.o(.text+0x69): undefined reference to `mysql_error'
/tmp/ccFKsdPa.o(.text+0x9a): undefined reference to `mysql_close'
```

You should be able to solve this problem by adding `-L`*`dir_path`* `-lmysqlclient` at the end of your link command, where *`dir_path`* represents the path name of the directory where the client library is located. To determine the correct directory, try this command:

```
shell> mysql_config --libs
```

The output from `mysql_config` might indicate other libraries that should be specified on the link command as well. You can include `mysql_config` output directly in your compile or link command using backticks. For example:

```
shell> gcc -o progname progname.o `mysql_config --libs`
```

If an error occurs at link time that the `floor` symbol is undefined, link to the math library by adding `-lm` to the end of the compile/link line. Similarly, if you get undefined-reference errors for other functions that should exist on your system, such as `connect()`, check the manual page for the function in question to determine which libraries you should add to the link command.

If you get undefined-reference errors such as the following for functions that do not exist on your system, it usually means that your MySQL client library was compiled on a system that is not 100% compatible with yours:

```
mf_format.o(.text+0x201): undefined reference to `__lxstat'
```

In this case, you should download the latest MySQL or Connector/C source distribution and compile the MySQL client library yourself. See Section 2.8, “Installing MySQL from Source”, and MySQL Connector/C Developer Guide.

27.7.4.2 Building C API Client Programs Using pkg-config

MySQL distributions contain a `mysqlclient.pc` file that provides information about MySQL configuration for use by the `pkg-config` command. This enables `pkg-config` to be used as an

alternative to `mysql_config` for obtaining information such as compiler flags or link libraries required to compile MySQL applications. For example, the following pairs of commands are equivalent:

```
mysql_config --cflags
pkg-config --cflags mysqlclient

mysql_config --libs
pkg-config --libs mysqlclient
```

The last `pkg-config` command produces flags for dynamic linking. To produce flags for static linking, use this command:

```
pkg-config --static --libs mysqlclient
```

On some platforms, the output with and without `--static` might be the same.

Note

If `pkg-config` does not find MySQL information, it might be necessary to set the `PKG_CONFIG_PATH` environment variable to the directory in which the `mysqlclient.pc` file is located, which by default is usually the `pkgconfig` directory under the MySQL library directory. For example (adjust the location appropriately):

```
export PKG_CONFIG_PATH=/usr/local/mysql/lib/pkgconfig # sh, bash, ...
setenv PKG_CONFIG_PATH /usr/local/mysql/lib/pkgconfig # csh, tcsh, ...
```

The `mysqlconfig.pc` installation location can be controlled using the `INSTALL_PKGCONFIGDIR CMake` option. See Section 2.8.4, “MySQL Source-Configuration Options”.

The `--variable` option takes a configuration variable name and displays the variable value:

```
pkg-config --variable=prefix mysqlclient     # installation prefix directory
pkg-config --variable=includedir mysqlclient # header file directory
pkg-config --variable=libdir mysqlclient     # library directory
```

To see which variable values `pkg-config` can display using the `--variable` option, use this command:

```
pkg-config --print-variables mysqlclient
```

You can use `pkg-config` within a command line using backticks to include the output that it produces for particular options. For example, to compile and link a MySQL client program, use `pkg-config` as follows:

```
gcc -c `pkg-config --cflags mysqlclient` progname.c
gcc -o progname progname.o `pkg-config --libs mysqlclient`
```

27.7.4.3 Writing C API Threaded Client Programs

The client library is almost thread-safe. The biggest problem is that the subroutines in `sql/net_serv.cc` that read from sockets are not interrupt-safe. This was done with the thought that you might want to have your own alarm that can break a long read to a server. If you install interrupt handlers for the `SIGPIPE` interrupt, socket handling should be thread-safe.

To avoid aborting the program when a connection terminates, MySQL blocks `SIGPIPE` on the first call to `mysql_library_init()`, `mysql_init()`, or `mysql_connect()`. To use your own `SIGPIPE` handler, first call `mysql_library_init()`, then install your handler.

If “undefined symbol” errors occur when linking against the `libmysqlclient` client library, in most cases this is because you have not included the thread libraries on the link/compile command.

The client library is thread-safe per connection. You can let two threads share the same connection with the following caveats:

- Multiple threads cannot send a query to the MySQL server at the same time on the same connection. In particular, you must ensure that between calls to `mysql_query()` and `mysql_store_result()` in one thread, no other thread uses the same connection. You must have a mutex lock around your pair of `mysql_query()` and `mysql_store_result()` calls. After `mysql_store_result()` returns, the lock can be released and other threads may query the same connection.

  If you use POSIX threads, you can use `pthread_mutex_lock()` and `pthread_mutex_unlock()` to establish and release a mutex lock.

- Many threads can access different result sets that are retrieved with `mysql_store_result()`.

- To use `mysql_use_result()`, you must ensure that no other thread is using the same connection until the result set is closed. However, it really is best for threaded clients that share the same connection to use `mysql_store_result()`.

You need to know the following if you have a thread that did not create the connection to the MySQL database but is calling MySQL functions:

When you call `mysql_init()`, MySQL creates a thread-specific variable for the thread that is used by the debug library (among other things). If you call a MySQL function before the thread has called `mysql_init()`, the thread does not have the necessary thread-specific variables in place and you are likely to end up with a core dump sooner or later. To avoid problems, you must do the following:

1. Call `mysql_library_init()` before any other MySQL functions. It is not thread-safe, so call it before threads are created, or protect the call with a mutex.

2. Arrange for `mysql_thread_init()` to be called early in the thread handler before calling any MySQL function. If you call `mysql_init()`, it will call `mysql_thread_init()` for you.

3. In the thread, call `mysql_thread_end()` before calling `pthread_exit()`. This frees the memory used by MySQL thread-specific variables.

The preceding notes regarding `mysql_init()` also apply to `mysql_connect()`, which calls `mysql_init()`.

27.7.4.4 Running C API Client Programs

If, after an upgrade, you experience problems with compiled client programs, such as `Commands out of sync` or unexpected core dumps, the programs were probably compiled using old header or library files. In this case, check the date of the `mysql.h` file and `libmysqlclient.a` library used for compilation to verify that they are from the new MySQL distribution. If not, recompile the programs with the new headers and libraries. Recompilation might also be necessary for programs compiled against the shared client library if the library major version number has changed (for example, from `libmysqlclient.so.17` to `libmysqlclient.so.18`).

The major client library version determines compatibility. (For example, for `libmysqlclient.so.18.1.0`, the major version is 18.) For this reason, the libraries shipped with newer versions of MySQL are drop-in replacements for older versions that have the same major number. As long as the major library version is the same, you can upgrade the library and old applications should continue to work with it.

Undefined-reference errors might occur at runtime when you try to execute a MySQL program. If these errors specify symbols that start with `mysql_` or indicate that the `libmysqlclient` library cannot be found, it means that your system cannot find the shared `libmysqlclient.so` library. The solution

to this problem is to tell your system to search for shared libraries in the directory where that library is located. Use whichever of the following methods is appropriate for your system:

- Add the path of the directory where `libmysqlclient.so` is located to the `LD_LIBRARY_PATH` or `LD_LIBRARY` environment variable.
- On OS X, add the path of the directory where `libmysqlclient.dylib` is located to the `DYLD_LIBRARY_PATH` environment variable.
- Copy the shared-library files (such as `libmysqlclient.so`) to some directory that is searched by your system, such as `/lib`, and update the shared library information by executing `ldconfig`. Be sure to copy all related files. A shared library might exist under several names, using symlinks to provide the alternate names.

27.7.4.5 C API Server and Client Library Versions

The string and numeric forms of the MySQL server version are available at compile time as the values of the `MYSQL_SERVER_VERSION` and `MYSQL_VERSION_ID` macros, and at runtime as the values of the `mysql_get_server_info()` and `mysql_get_server_version()` functions.

The MySQL client library version depends on the type of distribution that provides the library:

- For MySQL distributions, the client library version is the MySQL version. The string and numeric forms of this version are available at compile time as the values of the `MYSQL_SERVER_VERSION` and `MYSQL_VERSION_ID` macros, and at runtime as the values of the `mysql_get_client_info()` and `mysql_get_client_version()` functions.

  The `LIBMYSQL_VERSION` and `LIBMYSQL_VERSION_ID` macros have the same values as `MYSQL_SERVER_VERSION` and `MYSQL_VERSION_ID` and the two sets of macros can be used interchangeably.

- For Connector/C distributions, the client library version is the Connector/C version. The string and numeric forms of this version are available at compile time as the values of the `LIBMYSQL_VERSION` and `LIBMYSQL_VERSION_ID` macros, and at runtime as the values of the `mysql_get_client_info()` and `mysql_get_client_version()` functions.

  The `MYSQL_SERVER_VERSION` and `MYSQL_VERSION_ID` macros indicate the string and numeric forms of the MySQL version on which the Connector/C distribution is based.

27.7.5 C API Data Structures

This section describes C API data structures other than those used for prepared statements or the replication stream interface. For information about those, see Section 27.7.9, “C API Prepared Statement Data Structures”, and Section 27.7.15, “C API Binary Log Data Structures”.

- `MYSQL`

  This structure represents handler for one database connection. It is used for almost all MySQL functions. Do not try to make a copy of a `MYSQL` structure. There is no guarantee that such a copy will be usable.

- `MYSQL_RES`

  This structure represents the result of a query that returns rows (`SELECT`, `SHOW`, `DESCRIBE`, `EXPLAIN`). The information returned from a query is called the *result set* in the remainder of this section.

- `MYSQL_ROW`

  This is a type-safe representation of one row of data. It is currently implemented as an array of counted byte strings. (You cannot treat these as null-terminated strings if field values may contain

binary data, because such values may contain null bytes internally.) Rows are obtained by calling `mysql_fetch_row()`.

- `MYSQL_FIELD`

  This structure contains metadata: information about a field, such as the field's name, type, and size. Its members are described in more detail later in this section. You may obtain the `MYSQL_FIELD` structures for each field by calling `mysql_fetch_field()` repeatedly. Field values are not part of this structure; they are contained in a `MYSQL_ROW` structure.

- `MYSQL_FIELD_OFFSET`

  This is a type-safe representation of an offset into a MySQL field list. (Used by `mysql_field_seek()`.) Offsets are field numbers within a row, beginning at zero.

- `my_ulonglong`

  The type used for the number of rows and for `mysql_affected_rows()`, `mysql_num_rows()`, and `mysql_insert_id()`. This type provides a range of `0` to `1.84e19`.

  Some functions that return a row count using this type return -1 as an unsigned value to indicate an error or exceptional condition. You can check for -1 by comparing the return value to `(my_ulonglong)-1` (or to `(my_ulonglong)~0`, which is equivalent).

  On some systems, attempting to print a value of type `my_ulonglong` does not work. To print such a value, convert it to `unsigned long` and use a `%lu` print format. Example:

```
printf ("Number of rows: %lu\n",
        (unsigned long) mysql_num_rows(result));
```

- `my_bool`

  A boolean type, for values that are true (nonzero) or false (zero). The `my_bool` type was used before MySQL 8.0. As of MySQL 8.0, use the `bool` or `int` C type instead.

The `MYSQL_FIELD` structure contains the members described in the following list. The definitions apply primarily for columns of result sets such as those produced by `SELECT` statements. `MYSQL_FIELD` structures are also used to provide metadata for `OUT` and `INOUT` parameters returned from stored procedures executed using prepared `CALL` statements. For such parameters, some of the structure members have a meaning different from the meaning for column values.

- `char * name`

  The name of the field, as a null-terminated string. If the field was given an alias with an `AS` clause, the value of `name` is the alias. For a procedure parameter, the parameter name.

- `char * org_name`

  The name of the field, as a null-terminated string. Aliases are ignored. For expressions, the value is an empty string. For a procedure parameter, the parameter name.

- `char * table`

  The name of the table containing this field, if it is not a calculated field. For calculated fields, the `table` value is an empty string. If the column is selected from a view, `table` names the view. If the table or view was given an alias with an `AS` clause, the value of `table` is the alias. For a `UNION`, the value is the empty string. For a procedure parameter, the procedure name.

- `char * org_table`

  The name of the table, as a null-terminated string. Aliases are ignored. If the column is selected from a view, `org_table` names the view. If the column is selected from a derived table, `org_table`

names the base table. If a derived table wraps a view, `org_table` still names the base table. If the column is an expression, `org_table` is the empty string. For a `UNION`, the value is the empty string. For a procedure parameter, the value is the procedure name.

- `char * db`

  The name of the database that the field comes from, as a null-terminated string. If the field is a calculated field, `db` is an empty string. For a `UNION`, the value is the empty string. For a procedure parameter, the name of the database containing the procedure.

- `char * catalog`

  The catalog name. This value is always `"def"`.

- `char * def`

  The default value of this field, as a null-terminated string. This is set only if you use `mysql_list_fields()`.

- `unsigned long length`

  The width of the field. This corresponds to the display length, in bytes.

  The server determines the `length` value before it generates the result set, so this is the minimum length required for a data type capable of holding the largest possible value from the result column, without knowing in advance the actual values that will be produced by the query for the result set.

- `unsigned long max_length`

  The maximum width of the field for the result set (the length in bytes of the longest field value for the rows actually in the result set). If you use `mysql_store_result()` or `mysql_list_fields()`, this contains the maximum length for the field. If you use `mysql_use_result()`, the value of this variable is zero.

  The value of `max_length` is the length of the string representation of the values in the result set. For example, if you retrieve a `FLOAT` column and the "widest" value is `-12.345`, `max_length` is 7 (the length of `'-12.345'`).

  If you are using prepared statements, `max_length` is not set by default because for the binary protocol the lengths of the values depend on the types of the values in the result set. (See Section 27.7.9, "C API Prepared Statement Data Structures".) If you want the `max_length` values anyway, enable the `STMT_ATTR_UPDATE_MAX_LENGTH` option with `mysql_stmt_attr_set()` and the lengths will be set when you call `mysql_stmt_store_result()`. (See Section 27.7.11.3, "mysql_stmt_attr_set()", and Section 27.7.11.28, "mysql_stmt_store_result()".)

- `unsigned int name_length`

  The length of `name`.

- `unsigned int org_name_length`

  The length of `org_name`.

- `unsigned int table_length`

  The length of `table`.

- `unsigned int org_table_length`

  The length of `org_table`.

- `unsigned int db_length`

The length of `db`.

- `unsigned int catalog_length`

  The length of `catalog`.

- `unsigned int def_length`

  The length of `def`.

- `unsigned int flags`

  Bit-flags that describe the field. The `flags` value may have zero or more of the bits set that are shown in the following table.

Flag Value	Flag Description
`NOT_NULL_FLAG`	Field cannot be `NULL`
`PRI_KEY_FLAG`	Field is part of a primary key
`UNIQUE_KEY_FLAG`	Field is part of a unique key
`MULTIPLE_KEY_FLAG`	Field is part of a nonunique key
`UNSIGNED_FLAG`	Field has the `UNSIGNED` attribute
`ZEROFILL_FLAG`	Field has the `ZEROFILL` attribute
`BINARY_FLAG`	Field has the `BINARY` attribute
`AUTO_INCREMENT_FLAG`	Field has the `AUTO_INCREMENT` attribute
`ENUM_FLAG`	Field is an `ENUM`
`SET_FLAG`	Field is a `SET`
`BLOB_FLAG`	Field is a `BLOB` or `TEXT` (deprecated)
`TIMESTAMP_FLAG`	Field is a `TIMESTAMP` (deprecated)
`NUM_FLAG`	Field is numeric; see additional notes following table
`NO_DEFAULT_VALUE_FLAG`	Field has no default value; see additional notes following table

Some of these flags indicate data type information and are superseded by or used in conjunction with the `MYSQL_TYPE_xxx` value in the `field->type` member described later:

- To check for `BLOB` or `TIMESTAMP` values, check whether `type` is `MYSQL_TYPE_BLOB` or `MYSQL_TYPE_TIMESTAMP`. (The `BLOB_FLAG` and `TIMESTAMP_FLAG` flags are unneeded.)
- `ENUM` and `SET` values are returned as strings. For these, check that the `type` value is `MYSQL_TYPE_STRING` and that the `ENUM_FLAG` or `SET_FLAG` flag is set in the `flags` value.

`NUM_FLAG` indicates that a column is numeric. This includes columns with a type of `MYSQL_TYPE_DECIMAL`, `MYSQL_TYPE_NEWDECIMAL`, `MYSQL_TYPE_TINY`, `MYSQL_TYPE_SHORT`, `MYSQL_TYPE_LONG`, `MYSQL_TYPE_FLOAT`, `MYSQL_TYPE_DOUBLE`, `MYSQL_TYPE_NULL`, `MYSQL_TYPE_LONGLONG`, `MYSQL_TYPE_INT24`, and `MYSQL_TYPE_YEAR`.

`NO_DEFAULT_VALUE_FLAG` indicates that a column has no `DEFAULT` clause in its definition. This does not apply to `NULL` columns (because such columns have a default of `NULL`), or to `AUTO_INCREMENT` columns (which have an implied default value).

The following example illustrates a typical use of the `flags` value:

```
if (field->flags & NOT_NULL_FLAG)
    printf("Field cannot be null\n");
```

You may use the convenience macros shown in the following table to determine the boolean status of the `flags` value.

Flag Status	Description
`IS_NOT_NULL(flags)`	True if this field is defined as `NOT NULL`
`IS_PRI_KEY(flags)`	True if this field is a primary key
`IS_BLOB(flags)`	True if this field is a `BLOB` or `TEXT` (deprecated; test `field->type` instead)

- `unsigned int decimals`

  The number of decimals for numeric fields, and the fractional seconds precision for temporal fields.

- `unsigned int charsetnr`

  An ID number that indicates the character set/collation pair for the field.

  Normally, character values in result sets are converted to the character set indicated by the `character_set_results` system variable. In this case, `charsetnr` corresponds to the character set indicated by that variable. Character set conversion can be suppressed by setting `character_set_results` to `NULL`. In this case, `charsetnr` corresponds to the character set of the original table column or expression. See also Section 10.1.4, “Connection Character Sets and Collations”.

  To distinguish between binary and nonbinary data for string data types, check whether the `charsetnr` value is 63. If so, the character set is `binary`, which indicates binary rather than nonbinary data. This enables you to distinguish `BINARY` from `CHAR`, `VARBINARY` from `VARCHAR`, and the `BLOB` types from the `TEXT` types.

  `charsetnr` values are the same as those displayed in the `Id` column of the `SHOW COLLATION` statement or the `ID` column of the `INFORMATION_SCHEMA COLLATIONS` table. You can use those information sources to see which character set and collation specific `charsetnr` values indicate:

```
mysql> SHOW COLLATION WHERE Id = 63;
+----------+---------+----+---------+----------+---------+
| Collation | Charset | Id | Default | Compiled | Sortlen |
+----------+---------+----+---------+----------+---------+
| binary    | binary  | 63 | Yes     | Yes      |       1 |
+----------+---------+----+---------+----------+---------+

mysql> SELECT COLLATION_NAME, CHARACTER_SET_NAME
    -> FROM INFORMATION_SCHEMA.COLLATIONS WHERE ID = 33;
+----------------+--------------------+
| COLLATION_NAME  | CHARACTER_SET_NAME |
+----------------+--------------------+
| utf8_general_ci | utf8               |
+----------------+--------------------+
```

- `enum enum_field_types type`

  The type of the field. The `type` value may be one of the `MYSQL_TYPE_` symbols shown in the following table.

Type Value	Type Description
`MYSQL_TYPE_TINY`	`TINYINT` field
`MYSQL_TYPE_SHORT`	`SMALLINT` field
`MYSQL_TYPE_LONG`	`INTEGER` field
`MYSQL_TYPE_INT24`	`MEDIUMINT` field

Type Value	Type Description
MYSQL_TYPE_LONGLONG	BIGINT field
MYSQL_TYPE_DECIMAL	DECIMAL or NUMERIC field
MYSQL_TYPE_NEWDECIMAL	Precision math DECIMAL or NUMERIC
MYSQL_TYPE_FLOAT	FLOAT field
MYSQL_TYPE_DOUBLE	DOUBLE or REAL field
MYSQL_TYPE_BIT	BIT field
MYSQL_TYPE_TIMESTAMP	TIMESTAMP field
MYSQL_TYPE_DATE	DATE field
MYSQL_TYPE_TIME	TIME field
MYSQL_TYPE_DATETIME	DATETIME field
MYSQL_TYPE_YEAR	YEAR field
MYSQL_TYPE_STRING	CHAR or BINARY field
MYSQL_TYPE_VAR_STRING	VARCHAR or VARBINARY field
MYSQL_TYPE_BLOB	BLOB or TEXT field (use max_length to determine the maximum length)
MYSQL_TYPE_SET	SET field
MYSQL_TYPE_ENUM	ENUM field
MYSQL_TYPE_GEOMETRY	Spatial field
MYSQL_TYPE_NULL	NULL-type field

The MYSQL_TYPE_TIME2, MYSQL_TYPE_DATETIME2, and MYSQL_TYPE_TIMESTAMP2) type codes are used only on the server side. Clients see the MYSQL_TYPE_TIME, MYSQL_TYPE_DATETIME, and MYSQL_TYPE_TIMESTAMP codes.

You can use the IS_NUM() macro to test whether a field has a numeric type. Pass the type value to IS_NUM() and it evaluates to TRUE if the field is numeric:

```
if (IS_NUM(field->type))
    printf("Field is numeric\n");
```

ENUM and SET values are returned as strings. For these, check that the type value is MYSQL_TYPE_STRING and that the ENUM_FLAG or SET_FLAG flag is set in the flags value.

27.7.6 C API Function Overview

The following table summarizes the functions available in the C API. For greater detail, see the descriptions in Section 27.7.7, “C API Function Descriptions”.

Table 27.3 C API Function Names and Descriptions

Function	Description
my_init()	Initialize global variables, and thread handler in thread-safe programs
mysql_affected_rows()	Returns the number of rows changed/deleted/inserted by the last UPDATE, DELETE, or INSERT query
mysql_autocommit()	Toggles autocommit mode on/off
mysql_change_user()	Changes user and database on an open connection
mysql_character_set_name()	Return default character set name for current connection
mysql_client_find_plugin()	Return pointer to plugin

Function	Description
mysql_client_register_plugin()	Register a plugin
mysql_close()	Closes a server connection
mysql_commit()	Commits the transaction
mysql_connect()	Connects to a MySQL server (this function is deprecated; use mysql_real_connect() instead)
mysql_create_db()	Creates a database (this function is deprecated; use the SQL statement CREATE DATABASE instead)
mysql_data_seek()	Seeks to an arbitrary row number in a query result set
mysql_debug()	Does a DBUG_PUSH with the given string
mysql_drop_db()	Drops a database (this function is deprecated; use the SQL statement DROP DATABASE instead)
mysql_dump_debug_info()	Makes the server write debug information to the log
mysql_eof()	Determines whether the last row of a result set has been read (this function is deprecated; mysql_errno() or mysql_error() may be used instead)
mysql_errno()	Returns the error number for the most recently invoked MySQL function
mysql_error()	Returns the error message for the most recently invoked MySQL function
mysql_escape_string()	Escapes special characters in a string for use in an SQL statement
mysql_fetch_field()	Returns the type of the next table field
mysql_fetch_field_direct()	Returns the type of a table field, given a field number
mysql_fetch_fields()	Returns an array of all field structures
mysql_fetch_lengths()	Returns the lengths of all columns in the current row
mysql_fetch_row()	Fetches the next row from the result set
mysql_field_count()	Returns the number of result columns for the most recent statement
mysql_field_seek()	Puts the column cursor on a specified column
mysql_field_tell()	Returns the position of the field cursor used for the last mysql_fetch_field()
mysql_free_result()	Frees memory used by a result set
mysql_get_character_set_info()	Return information about default character set
mysql_get_client_info()	Returns client version information as a string
mysql_get_client_version()	Returns client version information as an integer
mysql_get_host_info()	Returns a string describing the connection
mysql_get_option()	Returns the value of a mysql_options() option
mysql_get_proto_info()	Returns the protocol version used by the connection
mysql_get_server_info()	Returns the server version number
mysql_get_server_version()	Returns version number of server as an integer
mysql_get_ssl_cipher()	Return current SSL cipher
mysql_hex_string()	Encode string in hexadecimal format
mysql_info()	Returns information about the most recently executed query
mysql_init()	Gets or initializes a MYSQL structure

Function	Description
mysql_insert_id()	Returns the ID generated for an AUTO_INCREMENT column by the previous query
mysql_kill()	Kills a given thread
mysql_library_end()	Finalize the MySQL C API library
mysql_library_init()	Initialize the MySQL C API library
mysql_list_dbs()	Returns database names matching a simple regular expression
mysql_list_fields()	Returns field names matching a simple regular expression
mysql_list_processes()	Returns a list of the current server threads
mysql_list_tables()	Returns table names matching a simple regular expression
mysql_load_plugin()	Load a plugin
mysql_load_plugin_v()	Load a plugin
mysql_more_results()	Checks whether any more results exist
mysql_next_result()	Returns/initiates the next result in multiple-result executions
mysql_num_fields()	Returns the number of columns in a result set
mysql_num_rows()	Returns the number of rows in a result set
mysql_options()	Sets connect options for mysql_real_connect()
mysql_options4()	Sets connect options for mysql_real_connect()
mysql_ping()	Checks whether the connection to the server is working, reconnecting as necessary
mysql_plugin_options()	Set a plugin option
mysql_query()	Executes an SQL query specified as a null-terminated string
mysql_real_connect()	Connects to a MySQL server
mysql_real_escape_string()	Escapes special characters in a string for use in an SQL statement, taking into account the current character set of the connection
mysql_real_escape_string_quote()	Escapes special characters in a string for use in an SQL statement, taking into account the current character set of the connection and the quoting context
mysql_real_query()	Executes an SQL query specified as a counted string
mysql_refresh()	Flush or reset tables and caches
mysql_reload()	Tells the server to reload the grant tables
mysql_reset_connection()	Reset connection to clear session state
mysql_rollback()	Rolls back the transaction
mysql_row_seek()	Seeks to a row offset in a result set, using value returned from mysql_row_tell()
mysql_row_tell()	Returns the row cursor position
mysql_select_db()	Selects a database
mysql_server_end()	Finalize the MySQL C API library
mysql_server_init()	Initialize the MySQL C API library
mysql_session_track_get_first()	Get first part of session state-change information
mysql_session_track_get_next()	Get next part of session state-change information
mysql_set_character_set()	Set default character set for current connection
mysql_set_local_infile_default()	Set the LOAD DATA LOCAL INFILE handler callbacks to their default values

Function	Description
mysql_set_local_infile_handler()	Install application-specific LOAD DATA LOCAL INFILE handler callbacks
mysql_set_server_option()	Sets an option for the connection (like multi-statements)
mysql_sqlstate()	Returns the SQLSTATE error code for the last error
mysql_shutdown()	Shuts down the database server
mysql_ssl_set()	Prepare to establish SSL connection to server
mysql_stat()	Returns the server status as a string
mysql_store_result()	Retrieves a complete result set to the client
mysql_thread_end()	Finalize thread handler
mysql_thread_id()	Returns the current thread ID
mysql_thread_init()	Initialize thread handler
mysql_thread_safe()	Returns 1 if the clients are compiled as thread-safe
mysql_use_result()	Initiates a row-by-row result set retrieval
mysql_warning_count()	Returns the warning count for the previous SQL statement

Application programs should use this general outline for interacting with MySQL:

1. Initialize the MySQL library by calling mysql_library_init().
2. Initialize a connection handler by calling mysql_init() and connect to the server by calling mysql_real_connect().
3. Issue SQL statements and process their results. (The following discussion provides more information about how to do this.)
4. Close the connection to the MySQL server by calling mysql_close().
5. End use of the MySQL library by calling mysql_library_end().

The purpose of calling mysql_library_init() and mysql_library_end() is to provide proper initialization and finalization of the MySQL library. For applications that are linked with the client library, they provide improved memory management. If you do not call mysql_library_end(), a block of memory remains allocated. (This does not increase the amount of memory used by the application, but some memory leak detectors will complain about it.) For applications that are linked with the embedded server, these calls start and stop the server.

In a nonmulti-threaded environment, the call to mysql_library_init() may be omitted, because mysql_init() will invoke it automatically as necessary. However, mysql_library_init() is not thread-safe in a multi-threaded environment, and thus neither is mysql_init(), which calls mysql_library_init(). You must either call mysql_library_init() prior to spawning any threads, or else use a mutex to protect the call, whether you invoke mysql_library_init() or indirectly through mysql_init(). This should be done prior to any other client library call.

To connect to the server, call mysql_init() to initialize a connection handler, then call mysql_real_connect() with that handler (along with other information such as the host name, user name, and password). Upon connection, mysql_real_connect() sets the reconnect flag (part of the MYSQL structure) to a value of 1 in versions of the API older than 5.0.3, or 0 in newer versions. A value of 1 for this flag indicates that if a statement cannot be performed because of a lost connection, to try reconnecting to the server before giving up. You can use the MYSQL_OPT_RECONNECT option to mysql_options() to control reconnection behavior. When you are done with the connection, call mysql_close() to terminate it.

While a connection is active, the client may send SQL statements to the server using mysql_query() or mysql_real_query(). The difference between the two is that mysql_query() expects the query

to be specified as a null-terminated string whereas `mysql_real_query()` expects a counted string. If the string contains binary data (which may include null bytes), you must use `mysql_real_query()`.

For each non-`SELECT` query (for example, `INSERT`, `UPDATE`, `DELETE`), you can find out how many rows were changed (affected) by calling `mysql_affected_rows()`.

For `SELECT` queries, you retrieve the selected rows as a result set. (Note that some statements are `SELECT`-like in that they return rows. These include `SHOW`, `DESCRIBE`, and `EXPLAIN`. Treat these statements the same way as `SELECT` statements.)

There are two ways for a client to process result sets. One way is to retrieve the entire result set all at once by calling `mysql_store_result()`. This function acquires from the server all the rows returned by the query and stores them in the client. The second way is for the client to initiate a row-by-row result set retrieval by calling `mysql_use_result()`. This function initializes the retrieval, but does not actually get any rows from the server.

In both cases, you access rows by calling `mysql_fetch_row()`. With `mysql_store_result()`, `mysql_fetch_row()` accesses rows that have previously been fetched from the server. With `mysql_use_result()`, `mysql_fetch_row()` actually retrieves the row from the server. Information about the size of the data in each row is available by calling `mysql_fetch_lengths()`.

After you are done with a result set, call `mysql_free_result()` to free the memory used for it.

The two retrieval mechanisms are complementary. Choose the approach that is most appropriate for each client application. In practice, clients tend to use `mysql_store_result()` more commonly.

An advantage of `mysql_store_result()` is that because the rows have all been fetched to the client, you not only can access rows sequentially, you can move back and forth in the result set using `mysql_data_seek()` or `mysql_row_seek()` to change the current row position within the result set. You can also find out how many rows there are by calling `mysql_num_rows()`. On the other hand, the memory requirements for `mysql_store_result()` may be very high for large result sets and you are more likely to encounter out-of-memory conditions.

An advantage of `mysql_use_result()` is that the client requires less memory for the result set because it maintains only one row at a time (and because there is less allocation overhead, `mysql_use_result()` can be faster). Disadvantages are that you must process each row quickly to avoid tying up the server, you do not have random access to rows within the result set (you can only access rows sequentially), and the number of rows in the result set is unknown until you have retrieved them all. Furthermore, you *must* retrieve all the rows even if you determine in mid-retrieval that you've found the information you were looking for.

The API makes it possible for clients to respond appropriately to statements (retrieving rows only as necessary) without knowing whether the statement is a `SELECT`. You can do this by calling `mysql_store_result()` after each `mysql_query()` (or `mysql_real_query()`). If the result set call succeeds, the statement was a `SELECT` and you can read the rows. If the result set call fails, call `mysql_field_count()` to determine whether a result was actually to be expected. If `mysql_field_count()` returns zero, the statement returned no data (indicating that it was an `INSERT`, `UPDATE`, `DELETE`, and so forth), and was not expected to return rows. If `mysql_field_count()` is nonzero, the statement should have returned rows, but did not. This indicates that the statement was a `SELECT` that failed. See the description for `mysql_field_count()` for an example of how this can be done.

Both `mysql_store_result()` and `mysql_use_result()` enable you to obtain information about the fields that make up the result set (the number of fields, their names and types, and so forth). You can access field information sequentially within the row by calling `mysql_fetch_field()` repeatedly, or by field number within the row by calling `mysql_fetch_field_direct()`. The current field cursor position may be changed by calling `mysql_field_seek()`. Setting the field cursor affects subsequent calls to `mysql_fetch_field()`. You can also get information for fields all at once by calling `mysql_fetch_fields()`.

For detecting and reporting errors, MySQL provides access to error information by means of the `mysql_errno()` and `mysql_error()` functions. These return the error code or error message for the most recently invoked function that can succeed or fail, enabling you to determine when an error occurred and what it was.

27.7.7 C API Function Descriptions

In the descriptions here, a parameter or return value of `NULL` means `NULL` in the sense of the C programming language, not a MySQL `NULL` value.

Functions that return a value generally return a pointer or an integer. Unless specified otherwise, functions returning a pointer return a non-`NULL` value to indicate success or a `NULL` value to indicate an error, and functions returning an integer return zero to indicate success or nonzero to indicate an error. Note that “nonzero” means just that. Unless the function description says otherwise, do not test against a value other than zero:

```
if (result)                  /* correct */
    ... error ...

if (result < 0)              /* incorrect */
    ... error ...

if (result == -1)            /* incorrect */
    ... error ...
```

When a function returns an error, the **Errors** subsection of the function description lists the possible types of errors. You can find out which of these occurred by calling `mysql_errno()`. A string representation of the error may be obtained by calling `mysql_error()`.

27.7.7.1 mysql_affected_rows()

```
my_ulonglong mysql_affected_rows(MYSQL *mysql)
```

Description

`mysql_affected_rows()` may be called immediately after executing a statement with `mysql_query()` or `mysql_real_query()`. It returns the number of rows changed, deleted, or inserted by the last statement if it was an `UPDATE`, `DELETE`, or `INSERT`. For `SELECT` statements, `mysql_affected_rows()` works like `mysql_num_rows()`.

For `UPDATE` statements, the affected-rows value by default is the number of rows actually changed. If you specify the `CLIENT_FOUND_ROWS` flag to `mysql_real_connect()` when connecting to `mysqld`, the affected-rows value is the number of rows “found”; that is, matched by the `WHERE` clause.

For `REPLACE` statements, the affected-rows value is 2 if the new row replaced an old row, because in this case, one row was inserted after the duplicate was deleted.

For `INSERT ... ON DUPLICATE KEY UPDATE` statements, the affected-rows value per row is 1 if the row is inserted as a new row, 2 if an existing row is updated, and 0 if an existing row is set to its current values. If you specify the `CLIENT_FOUND_ROWS` flag, the affected-rows value is 1 (not 0) if an existing row is set to its current values.

Following a `CALL` statement for a stored procedure, `mysql_affected_rows()` returns the value that it would return for the last statement executed within the procedure, or `0` if that statement would return `-1`. Within the procedure, you can use `ROW_COUNT()` at the SQL level to obtain the affected-rows value for individual statements.

`mysql_affected_rows()` returns a meaningful value for a wide range of statements. For details, see the description for `ROW_COUNT()` in Section 12.14, “Information Functions”.

Return Values

An integer greater than zero indicates the number of rows affected or retrieved. Zero indicates that no records were updated for an `UPDATE` statement, no rows matched the `WHERE` clause in the query or that no query has yet been executed. -1 indicates that the query returned an error or that, for a `SELECT` query, `mysql_affected_rows()` was called prior to calling `mysql_store_result()`.

Because `mysql_affected_rows()` returns an unsigned value, you can check for -1 by comparing the return value to `(my_ulonglong)-1` (or to `(my_ulonglong)~0`, which is equivalent).

Errors

None.

Example

```
char *stmt = "UPDATE products SET cost=cost*1.25
              WHERE group=10";
mysql_query(&mysql,stmt);
printf("%ld products updated",
       (long) mysql_affected_rows(&mysql));
```

27.7.7.2 mysql_autocommit()

```
bool mysql_autocommit(MYSQL *mysql, bool mode)
```

Description

Sets autocommit mode on if `mode` is 1, off if `mode` is 0.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

None.

27.7.7.3 mysql_change_user()

```
bool mysql_change_user(MYSQL *mysql, const char *user, const char *password,
const char *db)
```

Description

Changes the user and causes the database specified by `db` to become the default (current) database on the connection specified by `mysql`. In subsequent queries, this database is the default for table references that include no explicit database specifier.

`mysql_change_user()` fails if the connected user cannot be authenticated or does not have permission to use the database. In this case, the user and database are not changed.

Pass a `db` parameter of `NULL` if you do not want to have a default database.

This function resets the session state as if one had done a new connect and reauthenticated. (See Section 27.7.19, “Controlling Automatic Reconnection Behavior”.) It always performs a `ROLLBACK` of any active transactions, closes and drops all temporary tables, and unlocks all locked tables. Session system variables are reset to the values of the corresponding global system variables. Prepared statements are released and `HANDLER` variables are closed. Locks acquired with `GET_LOCK()` are released. These effects occur even if the user did not change.

To reset the connection state in a more lightweight manner without changing the user, use `mysql_reset_connection()`.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

The same that you can get from `mysql_real_connect()`, plus:

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

- `ER_UNKNOWN_COM_ERROR`

  The MySQL server does not implement this command (probably an old server).

- `ER_ACCESS_DENIED_ERROR`

  The user or password was wrong.

- `ER_BAD_DB_ERROR`

  The database did not exist.

- `ER_DBACCESS_DENIED_ERROR`

  The user did not have access rights to the database.

- `ER_WRONG_DB_NAME`

  The database name was too long.

Example

```
if (mysql_change_user(&mysql, "user", "password", "new_database"))
{
   fprintf(stderr, "Failed to change user.  Error: %s\n",
           mysql_error(&mysql));
}
```

27.7.7.4 mysql_character_set_name()

```
const char *mysql_character_set_name(MYSQL *mysql)
```

Description

Returns the default character set name for the current connection.

Return Values

The default character set name

Errors

None.

27.7.7.5 mysql_close()

```
void mysql_close(MYSQL *mysql)
```

Description

Closes a previously opened connection. `mysql_close()` also deallocates the connection handler pointed to by `mysql` if the handler was allocated automatically by `mysql_init()` or `mysql_connect()`.

Return Values

None.

Errors

None.

27.7.7.6 mysql_commit()

```
bool mysql_commit(MYSQL *mysql)
```

Description

Commits the current transaction.

The action of this function is subject to the value of the `completion_type` system variable. In particular, if the value of `completion_type` is `RELEASE` (or 2), the server performs a release after terminating a transaction and closes the client connection. Call `mysql_close()` from the client program to close the connection from the client side.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

None.

27.7.7.7 mysql_connect()

```
MYSQL *mysql_connect(MYSQL *mysql, const char *host, const char *user, const
char *passwd)
```

Description

This function is deprecated. Use `mysql_real_connect()` instead.

`mysql_connect()` attempts to establish a connection to a MySQL database engine running on `host`. `mysql_connect()` must complete successfully before you can execute any of the other API functions, with the exception of `mysql_get_client_info()`.

The meanings of the parameters are the same as for the corresponding parameters for `mysql_real_connect()` with the difference that the connection parameter may be `NULL`. In this case, the C API allocates memory for the connection structure automatically and frees it when you call `mysql_close()`. The disadvantage of this approach is that you cannot retrieve an error message if the connection fails. (To get error information from `mysql_errno()` or `mysql_error()`, you must provide a valid `MYSQL` pointer.)

Return Values

Same as for `mysql_real_connect()`.

Errors

Same as for `mysql_real_connect()`.

27.7.7.8 mysql_create_db()

```
int mysql_create_db(MYSQL *mysql, const char *db)
```

Description

Creates the database named by the `db` parameter.

This function is deprecated. It is preferable to use `mysql_query()` to issue an SQL `CREATE DATABASE` statement instead.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

```
if(mysql_create_db(&mysql, "my_database"))
{
   fprintf(stderr, "Failed to create new database.  Error: %s\n",
           mysql_error(&mysql));
}
```

27.7.7.9 mysql_data_seek()

```
void mysql_data_seek(MYSQL_RES *result, my_ulonglong offset)
```

Description

Seeks to an arbitrary row in a query result set. The `offset` value is a row number. Specify a value in the range from `0` to `mysql_num_rows(result)-1`.

This function requires that the result set structure contains the entire result of the query, so `mysql_data_seek()` may be used only in conjunction with `mysql_store_result()`, not with `mysql_use_result()`.

Return Values

None.

Errors

None.

27.7.7.10 mysql_debug()

```
void mysql_debug(const char *debug)
```

Description

Does a `DBUG_PUSH` with the given string. `mysql_debug()` uses the Fred Fish debug library. To use this function, you must compile the client library to support debugging. See Section 28.5.3, “The DBUG Package”.

Return Values

None.

Errors

None.

Example

The call shown here causes the client library to generate a trace file in `/tmp/client.trace` on the client machine:

```
mysql_debug("d:t:O,/tmp/client.trace");
```

27.7.7.11 mysql_drop_db()

```
int mysql_drop_db(MYSQL *mysql, const char *db)
```

Description

Drops the database named by the `db` parameter.

This function is deprecated. It is preferable to use `mysql_query()` to issue an SQL `DROP DATABASE` statement instead.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

```
if(mysql_drop_db(&mysql, "my_database"))
  fprintf(stderr, "Failed to drop the database: Error: %s\n",
          mysql_error(&mysql));
```

27.7.7.12 mysql_dump_debug_info()

```
int mysql_dump_debug_info(MYSQL *mysql)
```

Description

Instructs the server to write debugging information to the error log. The connected user must have the `SUPER` privilege.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.13 mysql_eof()

```
bool mysql_eof(MYSQL_RES *result)
```

Description

This function is deprecated. `mysql_errno()` or `mysql_error()` may be used instead.

`mysql_eof()` determines whether the last row of a result set has been read.

If you acquire a result set from a successful call to `mysql_store_result()`, the client receives the entire set in one operation. In this case, a `NULL` return from `mysql_fetch_row()` always means the end of the result set has been reached and it is unnecessary to call `mysql_eof()`. When used with `mysql_store_result()`, `mysql_eof()` always returns true.

On the other hand, if you use `mysql_use_result()` to initiate a result set retrieval, the rows of the set are obtained from the server one by one as you call `mysql_fetch_row()` repeatedly. Because an error may occur on the connection during this process, a `NULL` return value from `mysql_fetch_row()` does not necessarily mean the end of the result set was reached normally. In this case, you can use `mysql_eof()` to determine what happened. `mysql_eof()` returns a nonzero value if the end of the result set was reached and zero if an error occurred.

Historically, `mysql_eof()` predates the standard MySQL error functions `mysql_errno()` and `mysql_error()`. Because those error functions provide the same information, their use is preferred over `mysql_eof()`, which is deprecated. (In fact, they provide more information, because `mysql_eof()` returns only a boolean value whereas the error functions indicate a reason for the error when one occurs.)

Return Values

Zero for success. Nonzero if the end of the result set has been reached.

Errors

None.

Example

The following example shows how you might use `mysql_eof()`:

```
mysql_query(&mysql,"SELECT * FROM some_table");
result = mysql_use_result(&mysql);
while((row = mysql_fetch_row(result)))
{
    // do something with data
}
if(!mysql_eof(result))  // mysql_fetch_row() failed due to an error
{
    fprintf(stderr, "Error: %s\n", mysql_error(&mysql));
}
```

However, you can achieve the same effect with the standard MySQL error functions:

```
mysql_query(&mysql,"SELECT * FROM some_table");
result = mysql_use_result(&mysql);
while((row = mysql_fetch_row(result)))
{
    // do something with data
}
if(mysql_errno(&mysql))  // mysql_fetch_row() failed due to an error
{
    fprintf(stderr, "Error: %s\n", mysql_error(&mysql));
}
```

27.7.7.14 mysql_errno()

```
unsigned int mysql_errno(MYSQL *mysql)
```

Description

For the connection specified by `mysql`, `mysql_errno()` returns the error code for the most recently invoked API function that can succeed or fail. A return value of zero means that no error occurred. Client error message numbers are listed in the MySQL `errmsg.h` header file. Server error message numbers are listed in `mysqld_error.h`. Errors also are listed at Appendix B, *Errors, Error Codes, and Common Problems*.

Note

Some functions such as `mysql_fetch_row()` do not set `mysql_errno()` if they succeed. A rule of thumb is that all functions that have to ask the server for information reset `mysql_errno()` if they succeed.

MySQL-specific error numbers returned by `mysql_errno()` differ from SQLSTATE values returned by `mysql_sqlstate()`. For example, the `mysql` client program displays errors using

the following format, where `1146` is the `mysql_errno()` value and `'42S02'` is the corresponding `mysql_sqlstate()` value:

```
shell> SELECT * FROM no_such_table;
ERROR 1146 (42S02): Table 'test.no_such_table' doesn't exist
```

Return Values

An error code value for the last `mysql_xxx()` call, if it failed. zero means no error occurred.

Errors

None.

27.7.7.15 mysql_error()

```
const char *mysql_error(MYSQL *mysql)
```

Description

For the connection specified by `mysql`, `mysql_error()` returns a null-terminated string containing the error message for the most recently invoked API function that failed. If a function did not fail, the return value of `mysql_error()` may be the previous error or an empty string to indicate no error.

A rule of thumb is that all functions that have to ask the server for information reset `mysql_error()` if they succeed.

For functions that reset `mysql_error()`, either of these two tests can be used to check for an error:

```
if(*mysql_error(&mysql))
{
  // an error occurred
}

if(mysql_error(&mysql)[0])
{
  // an error occurred
}
```

The language of the client error messages may be changed by recompiling the MySQL client library. You can choose error messages in several different languages. See Section 10.2, “Setting the Error Message Language”.

Return Values

A null-terminated character string that describes the error. An empty string if no error occurred.

Errors

None.

27.7.7.16 mysql_escape_string()

Note

Do not use this function. `mysql_escape_string()` does not have arguments that enable it to respect the current character set or the quoting context. Use `mysql_real_escape_string_quote()` instead.

27.7.7.17 mysql_fetch_field()

```
MYSQL_FIELD *mysql_fetch_field(MYSQL_RES *result)
```

Description

Returns the definition of one column of a result set as a `MYSQL_FIELD` structure. Call this function repeatedly to retrieve information about all columns in the result set. `mysql_fetch_field()` returns `NULL` when no more fields are left.

`mysql_fetch_field()` is reset to return information about the first field each time you execute a new `SELECT` query. The field returned by `mysql_fetch_field()` is also affected by calls to `mysql_field_seek()`.

If you've called `mysql_query()` to perform a `SELECT` on a table but have not called `mysql_store_result()`, MySQL returns the default blob length (8KB) if you call `mysql_fetch_field()` to ask for the length of a `BLOB` field. (The 8KB size is chosen because MySQL does not know the maximum length for the `BLOB`. This should be made configurable sometime.) Once you've retrieved the result set, `field->max_length` contains the length of the largest value for this column in the specific query.

Return Values

The `MYSQL_FIELD` structure for the current column. `NULL` if no columns are left.

Errors

None.

Example

```
MYSQL_FIELD *field;

while((field = mysql_fetch_field(result)))
{
    printf("field name %s\n", field->name);
}
```

27.7.7.18 mysql_fetch_field_direct()

```
MYSQL_FIELD *mysql_fetch_field_direct(MYSQL_RES *result, unsigned int
fieldnr)
```

Description

Given a field number `fieldnr` for a column within a result set, returns that column's field definition as a `MYSQL_FIELD` structure. Use this function to retrieve the definition for an arbitrary column. Specify a value for `fieldnr` in the range from 0 to `mysql_num_fields(result)-1`.

Return Values

The `MYSQL_FIELD` structure for the specified column.

Errors

None.

Example

```
unsigned int num_fields;
unsigned int i;
MYSQL_FIELD *field;

num_fields = mysql_num_fields(result);
for(i = 0; i < num_fields; i++)
{
    field = mysql_fetch_field_direct(result, i);
```

```
    printf("Field %u is %s\n", i, field->name);
}
```

27.7.7.19 mysql_fetch_fields()

```
MYSQL_FIELD *mysql_fetch_fields(MYSQL_RES *result)
```

Description

Returns an array of all `MYSQL_FIELD` structures for a result set. Each structure provides the field definition for one column of the result set.

Return Values

An array of `MYSQL_FIELD` structures for all columns of a result set.

Errors

None.

Example

```
unsigned int num_fields;
unsigned int i;
MYSQL_FIELD *fields;

num_fields = mysql_num_fields(result);
fields = mysql_fetch_fields(result);
for(i = 0; i < num_fields; i++)
{
   printf("Field %u is %s\n", i, fields[i].name);
}
```

27.7.7.20 mysql_fetch_lengths()

```
unsigned long *mysql_fetch_lengths(MYSQL_RES *result)
```

Description

Returns the lengths of the columns of the current row within a result set. If you plan to copy field values, this length information is also useful for optimization, because you can avoid calling `strlen()`. In addition, if the result set contains binary data, you **must** use this function to determine the size of the data, because `strlen()` returns incorrect results for any field containing null characters.

The length for empty columns and for columns containing `NULL` values is zero. To see how to distinguish these two cases, see the description for `mysql_fetch_row()`.

Return Values

An array of unsigned long integers representing the size of each column (not including any terminating null bytes). `NULL` if an error occurred.

Errors

`mysql_fetch_lengths()` is valid only for the current row of the result set. It returns `NULL` if you call it before calling `mysql_fetch_row()` or after retrieving all rows in the result.

Example

```
MYSQL_ROW row;
unsigned long *lengths;
unsigned int num_fields;
unsigned int i;
```

```
row = mysql_fetch_row(result);
if (row)
{
    num_fields = mysql_num_fields(result);
    lengths = mysql_fetch_lengths(result);
    for(i = 0; i < num_fields; i++)
    {
         printf("Column %u is %lu bytes in length.\n",
                i, lengths[i]);
    }
}
```

27.7.7.21 mysql_fetch_row()

```
MYSQL_ROW mysql_fetch_row(MYSQL_RES *result)
```

Description

Retrieves the next row of a result set. When used after `mysql_store_result()`, `mysql_fetch_row()` returns `NULL` when there are no more rows to retrieve. When used after `mysql_use_result()`, `mysql_fetch_row()` returns `NULL` when there are no more rows to retrieve or if an error occurred.

The number of values in the row is given by `mysql_num_fields(result)`. If `row` holds the return value from a call to `mysql_fetch_row()`, pointers to the values are accessed as `row[0]` to `row[mysql_num_fields(result)-1]`. `NULL` values in the row are indicated by `NULL` pointers.

The lengths of the field values in the row may be obtained by calling `mysql_fetch_lengths()`. Empty fields and fields containing `NULL` both have length 0; you can distinguish these by checking the pointer for the field value. If the pointer is `NULL`, the field is `NULL`; otherwise, the field is empty.

Return Values

A `MYSQL_ROW` structure for the next row. `NULL` if there are no more rows to retrieve or if an error occurred.

Errors

Errors are not reset between calls to `mysql_fetch_row()`

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

```
MYSQL_ROW row;
unsigned int num_fields;
unsigned int i;

num_fields = mysql_num_fields(result);
while ((row = mysql_fetch_row(result)))
{
   unsigned long *lengths;
   lengths = mysql_fetch_lengths(result);
   for(i = 0; i < num_fields; i++)
   {
       printf("[%.*s] ", (int) lengths[i],
              row[i] ? row[i] : "NULL");
   }
```

```
   printf("\n");
}
```

27.7.7.22 mysql_field_count()

```
unsigned int mysql_field_count(MYSQL *mysql)
```

Description

Returns the number of columns for the most recent query on the connection.

The normal use of this function is when `mysql_store_result()` returned `NULL` (and thus you have no result set pointer). In this case, you can call `mysql_field_count()` to determine whether `mysql_store_result()` should have produced a nonempty result. This enables the client program to take proper action without knowing whether the query was a `SELECT` (or `SELECT`-like) statement. The example shown here illustrates how this may be done.

See Section 27.7.18.1, “Why mysql_store_result() Sometimes Returns NULL After mysql_query() Returns Success”.

Return Values

An unsigned integer representing the number of columns in a result set.

Errors

None.

Example

```
MYSQL_RES *result;
unsigned int num_fields;
unsigned int num_rows;

if (mysql_query(&mysql,query_string))
{
    // error
}
else // query succeeded, process any data returned by it
{
    result = mysql_store_result(&mysql);
    if (result)  // there are rows
    {
        num_fields = mysql_num_fields(result);
        // retrieve rows, then call mysql_free_result(result)
    }
    else  // mysql_store_result() returned nothing; should it have?
    {
        if(mysql_field_count(&mysql) == 0)
        {
            // query does not return data
            // (it was not a SELECT)
            num_rows = mysql_affected_rows(&mysql);
        }
        else // mysql_store_result() should have returned data
        {
            fprintf(stderr, "Error: %s\n", mysql_error(&mysql));
        }
    }
}
```

An alternative is to replace the `mysql_field_count(&mysql)` call with `mysql_errno(&mysql)`. In this case, you are checking directly for an error from `mysql_store_result()` rather than inferring from the value of `mysql_field_count()` whether the statement was a `SELECT`.

27.7.7.23 mysql_field_seek()

```
MYSQL_FIELD_OFFSET mysql_field_seek(MYSQL_RES *result, MYSQL_FIELD_OFFSET
offset)
```

Description

Sets the field cursor to the given offset. The next call to `mysql_fetch_field()` retrieves the field definition of the column associated with that offset.

To seek to the beginning of a row, pass an `offset` value of zero.

Return Values

The previous value of the field cursor.

Errors

None.

27.7.7.24 mysql_field_tell()

```
MYSQL_FIELD_OFFSET mysql_field_tell(MYSQL_RES *result)
```

Description

Returns the position of the field cursor used for the last `mysql_fetch_field()`. This value can be used as an argument to `mysql_field_seek()`.

Return Values

The current offset of the field cursor.

Errors

None.

27.7.7.25 mysql_free_result()

```
void mysql_free_result(MYSQL_RES *result)
```

Description

Frees the memory allocated for a result set by `mysql_store_result()`, `mysql_use_result()`, `mysql_list_dbs()`, and so forth. When you are done with a result set, you must free the memory it uses by calling `mysql_free_result()`.

Do not attempt to access a result set after freeing it.

Return Values

None.

Errors

None.

27.7.7.26 mysql_get_character_set_info()

```
void mysql_get_character_set_info(MYSQL *mysql, MY_CHARSET_INFO *cs)
```

Description

This function provides information about the default client character set. The default character set may be changed with the `mysql_set_character_set()` function.

Example

This example shows the fields that are available in the `MY_CHARSET_INFO` structure:

```
if (!mysql_set_character_set(&mysql, "utf8"))
{
    MY_CHARSET_INFO cs;
    mysql_get_character_set_info(&mysql, &cs);
    printf("character set information:\n");
    printf("character set+collation number: %d\n", cs.number);
    printf("character set name: %s\n", cs.name);
    printf("collation name: %s\n", cs.csname);
    printf("comment: %s\n", cs.comment);
    printf("directory: %s\n", cs.dir);
    printf("multi byte character min. length: %d\n", cs.mbminlen);
    printf("multi byte character max. length: %d\n", cs.mbmaxlen);
}
```

27.7.7.27 mysql_get_client_info()

```
const char *mysql_get_client_info(void)
```

Description

Returns a string that represents the MySQL client library version; for example, `"8.0.3"`.

The function value is the version of MySQL or Connector/C that provides the client library. For more information, see Section 27.7.4.5, “C API Server and Client Library Versions”.

Return Values

A character string that represents the MySQL client library version.

Errors

None.

27.7.7.28 mysql_get_client_version()

```
unsigned long mysql_get_client_version(void)
```

Description

Returns an integer that represents the MySQL client library version. The value has the format `XYYZZ` where `X` is the major version, `YY` is the release level (or minor version), and `ZZ` is the sub-version within the release level:

```
major_version*10000 + release_level*100 + sub_version
```

For example, `"8.0.3"` is returned as `80003`.

The function value is the version of MySQL or Connector/C that provides the client library. For more information, see Section 27.7.4.5, “C API Server and Client Library Versions”.

Return Values

An integer that represents the MySQL client library version.

Errors

None.

27.7.7.29 mysql_get_host_info()

```
const char *mysql_get_host_info(MYSQL *mysql)
```

Description

Returns a string describing the type of connection in use, including the server host name.

Return Values

A character string representing the server host name and the connection type.

Errors

None.

27.7.7.30 mysql_get_option()

```
int mysql_get_option(MYSQL *mysql, enum mysql_option option, const void
*arg)
```

Description

Returns the current value of an option settable using `mysql_options()`. The value should be treated as read only.

The `option` argument is the option for which you want its value. The `arg` argument is a pointer to a variable in which to store the option value. `arg` must be a pointer to a variable of the type appropriate for the `option` argument. The following table shows which variable type to use for each `option` value.

`arg` Type	**Applicable `option` Values**
`unsigned int`	`MYSQL_OPT_CONNECT_TIMEOUT`, `MYSQL_OPT_PROTOCOL`, `MYSQL_OPT_READ_TIMEOUT`, `MYSQL_OPT_SSL_MODE`, `MYSQL_OPT_RETRY_COUNT`, `MYSQL_OPT_WRITE_TIMEOUT`
`unsigned long`	`MYSQL_OPT_MAX_ALLOWED_PACKET`, `MYSQL_OPT_NET_BUFFER_LENGTH`
`bool`	`MYSQL_ENABLE_CLEARTEXT_PLUGIN`, `MYSQL_OPT_CAN_HANDLE_EXPIRED_PASSWORDS`, `MYSQL_OPT_COMPRESS`, `MYSQL_OPT_GUESS_CONNECTION`, `MYSQL_OPT_LOCAL_INFILE`, `MYSQL_OPT_RECONNECT`, `MYSQL_OPT_USE_EMBEDDED_CONNECTION`, `MYSQL_OPT_USE_REMOTE_CONNECTION`, `MYSQL_REPORT_DATA_TRUNCATION`, `MYSQL_SECURE_AUTH`
`const char *`	`MYSQL_DEFAULT_AUTH`, `MYSQL_OPT_BIND`, `MYSQL_OPT_SSL_CA`, `MYSQL_OPT_SSL_CAPATH` , `MYSQL_OPT_SSL_CERT`, `MYSQL_OPT_SSL_CIPHER`, `MYSQL_OPT_SSL_CRL`, `MYSQL_OPT_SSL_CRLPATH`, `MYSQL_OPT_SSL_KEY`, `MYSQL_OPT_TLS_VERSION`, `MYSQL_PLUGIN_DIR`, `MYSQL_READ_DEFAULT_FILE`, `MYSQL_READ_DEFAULT_GROUP`, `MYSQL_SERVER_PUBLIC_KEY`, `MYSQL_SET_CHARSET_DIR`, `MYSQL_SET_CHARSET_NAME`, `MYSQL_SET_CLIENT_IP`, `MYSQL_SHARED_MEMORY_BASE_NAME`
cannot be queried (error is returned)	`MYSQL_INIT_COMMAND`, `MYSQL_OPT_CONNECT_ATTR_DELETE`, `MYSQL_OPT_CONNECT_ATTR_RESET`, `MYSQL_OPT_NAMED_PIPE`

Return Values

Zero for success. Nonzero if an error occurred; this occurs for `option` values that cannot be queried.

Example

The following call tests the `MYSQL_OPT_RECONNECT` option. After the call returns successfully, the value of `reconnect` is true or false to indicate whether automatic reconnection is enabled.

```
bool reconnect;

if (mysql_get_option(mysql, MYSQL_OPT_RECONNECT, &reconnect))
  fprintf(stderr, "mysql_get_options() failed\n");
```

27.7.7.31 mysql_get_proto_info()

```
unsigned int mysql_get_proto_info(MYSQL *mysql)
```

Description

Returns the protocol version used by current connection.

Return Values

An unsigned integer representing the protocol version used by the current connection.

Errors

None.

27.7.7.32 mysql_get_server_info()

```
const char *mysql_get_server_info(MYSQL *mysql)
```

Description

Returns a string that represents the MySQL server version; for example, `"8.0.3"`.

Return Values

A character string that represents the MySQL server version.

Errors

None.

27.7.7.33 mysql_get_server_version()

```
unsigned long mysql_get_server_version(MYSQL *mysql)
```

Description

Returns an integer that represents the MySQL server version. The value has the format `XYYZZ` where `X` is the major version, `YY` is the release level (or minor version), and `ZZ` is the sub-version within the release level:

```
major_version*10000 + release_level*100 + sub_version
```

For example, `"8.0.3"` is returned as `80003`.

This function is useful in client programs for determining whether some version-specific server capability exists.

Return Values

An integer that represents the MySQL server version.

Errors

None.

27.7.7.34 mysql_get_ssl_cipher()

```
const char *mysql_get_ssl_cipher(MYSQL *mysql)
```

Description

`mysql_get_ssl_cipher()` returns the encryption cipher used for the given connection to the server. `mysql` is the connection handler returned from `mysql_init()`.

Return Values

A string naming the encryption cipher used for the connection, or `NULL` if no cipher is being used.

27.7.7.35 mysql_hex_string()

```
unsigned long mysql_hex_string(char *to, const char *from, unsigned long
length)
```

Description

This function creates a legal SQL string for use in an SQL statement. See Section 9.1.1, “String Literals”.

The string in the `from` argument is encoded in hexadecimal format, with each character encoded as two hexadecimal digits. The result is placed in the `to` argument, followed by a terminating null byte.

The string pointed to by `from` must be `length` bytes long. You must allocate the `to` buffer to be at least `length*2+1` bytes long. When `mysql_hex_string()` returns, the contents of `to` is a null-terminated string. The return value is the length of the encoded string, not including the terminating null byte.

The return value can be placed into an SQL statement using either `X'value'` or `0xvalue` format. However, the return value does not include the `X'...'` or `0x`. The caller must supply whichever of those is desired.

Example

```
char query[1000],*end;

end = strmov(query,"INSERT INTO test_table values(");
end = strmov(end,"X'");
end += mysql_hex_string(end,"What is this",12);
end = strmov(end,"',X'");
end += mysql_hex_string(end,"binary data: \0\r\n",16);
end = strmov(end,"')");

if (mysql_real_query(&mysql,query,(unsigned int) (end - query)))
{
   fprintf(stderr, "Failed to insert row, Error: %s\n",
           mysql_error(&mysql));
}
```

The `strmov()` function used in the example is included in the `libmysqlclient` library and works like `strcpy()` but returns a pointer to the terminating null of the first parameter.

Return Values

The length of the encoded string that is placed into `to`, not including the terminating null character.

Errors

None.

27.7.7.36 mysql_info()

```
const char *mysql_info(MYSQL *mysql)
```

Description

Retrieves a string providing information about the most recently executed statement, but only for the statements listed here. For other statements, `mysql_info()` returns `NULL`. The format of the string varies depending on the type of statement, as described here. The numbers are illustrative only; the string contains values appropriate for the statement.

- `INSERT INTO ... SELECT ...`

  String format: `Records: 100 Duplicates: 0 Warnings: 0`

- `INSERT INTO ... VALUES (...),(...),(...)...`

  String format: `Records: 3 Duplicates: 0 Warnings: 0`

- `LOAD DATA INFILE ...`

  String format: `Records: 1 Deleted: 0 Skipped: 0 Warnings: 0`

- `ALTER TABLE`

  String format: `Records: 3 Duplicates: 0 Warnings: 0`

- `UPDATE`

  String format: `Rows matched: 40 Changed: 40 Warnings: 0`

`mysql_info()` returns a non-`NULL` value for `INSERT ... VALUES` only for the multiple-row form of the statement (that is, only if multiple value lists are specified).

Return Values

A character string representing additional information about the most recently executed statement. `NULL` if no information is available for the statement.

Errors

None.

27.7.7.37 mysql_init()

```
MYSQL *mysql_init(MYSQL *mysql)
```

Description

Allocates or initializes a `MYSQL` object suitable for `mysql_real_connect()`. If `mysql` is a `NULL` pointer, the function allocates, initializes, and returns a new object. Otherwise, the object is initialized and the address of the object is returned. If `mysql_init()` allocates a new object, it is freed when `mysql_close()` is called to close the connection.

In a nonmulti-threaded environment, `mysql_init()` invokes `mysql_library_init()` automatically as necessary. However, `mysql_library_init()` is not thread-safe in a multi-threaded environment, and thus neither is `mysql_init()`. Before calling `mysql_init()`, either call `mysql_library_init()` prior to spawning any threads, or use a mutex to protect the `mysql_library_init()` call. This should be done prior to any other client library call.

Return Values

An initialized `MYSQL*` handler. `NULL` if there was insufficient memory to allocate a new object.

Errors

In case of insufficient memory, `NULL` is returned.

27.7.7.38 mysql_insert_id()

```
my_ulonglong mysql_insert_id(MYSQL *mysql)
```

Description

Returns the value generated for an AUTO_INCREMENT column by the previous INSERT or UPDATE statement. Use this function after you have performed an INSERT statement into a table that contains an AUTO_INCREMENT field, or have used INSERT or UPDATE to set a column value with LAST_INSERT_ID(*expr*).

The return value of mysql_insert_id() is always zero unless explicitly updated under one of the following conditions:

- INSERT statements that store a value into an AUTO_INCREMENT column. This is true whether the value is automatically generated by storing the special values NULL or 0 into the column, or is an explicit nonspecial value.
- In the case of a multiple-row INSERT statement, mysql_insert_id() returns the first automatically generated AUTO_INCREMENT value that was successfully inserted.

  If no rows are successfully inserted, mysql_insert_id() returns 0.
- If an INSERT ... SELECT statement is executed, and no automatically generated value is successfully inserted, mysql_insert_id() returns the ID of the last inserted row.
- If an INSERT ... SELECT statement uses LAST_INSERT_ID(*expr*), mysql_insert_id() returns *expr*.
- INSERT statements that generate an AUTO_INCREMENT value by inserting LAST_INSERT_ID(*expr*) into any column or by updating any column to LAST_INSERT_ID(*expr*).
- If the previous statement returned an error, the value of mysql_insert_id() is undefined.

The return value of mysql_insert_id() can be simplified to the following sequence:

1. If there is an AUTO_INCREMENT column, and an automatically generated value was successfully inserted, return the first such value.
2. If LAST_INSERT_ID(*expr*) occurred in the statement, return *expr*, even if there was an AUTO_INCREMENT column in the affected table.
3. The return value varies depending on the statement used. When called after an INSERT statement:
   - If there is an AUTO_INCREMENT column in the table, and there were some explicit values for this column that were successfully inserted into the table, return the last of the explicit values.

   When called after an INSERT ... ON DUPLICATE KEY UPDATE statement:
   - If there is an AUTO_INCREMENT column in the table and there were some explicit successfully inserted values or some updated values, return the last of the inserted or updated values.

mysql_insert_id() returns 0 if the previous statement does not use an AUTO_INCREMENT value. If you need to save the value for later, be sure to call mysql_insert_id() immediately after the statement that generates the value.

The value of mysql_insert_id() is affected only by statements issued within the current client connection. It is not affected by statements issued by other clients.

The LAST_INSERT_ID() SQL function will contain the value of the first automatically generated value that was successfully inserted. LAST_INSERT_ID() is not reset between statements because the

value of that function is maintained in the server. Another difference from `mysql_insert_id()` is that `LAST_INSERT_ID()` is not updated if you set an `AUTO_INCREMENT` column to a specific nonspecial value. See Section 12.14, “Information Functions”.

`mysql_insert_id()` returns `0` following a `CALL` statement for a stored procedure that generates an `AUTO_INCREMENT` value because in this case `mysql_insert_id()` applies to `CALL` and not the statement within the procedure. Within the procedure, you can use `LAST_INSERT_ID()` at the SQL level to obtain the `AUTO_INCREMENT` value.

The reason for the differences between `LAST_INSERT_ID()` and `mysql_insert_id()` is that `LAST_INSERT_ID()` is made easy to use in scripts while `mysql_insert_id()` tries to provide more exact information about what happens to the `AUTO_INCREMENT` column.

Return Values

Described in the preceding discussion.

Errors

- `ER_AUTO_INCREMENT_CONFLICT`

  A user-specified `AUTO_INCREMENT` value in a multi `INSERT` statement falls within the range between the current `AUTO_INCREMENT` value and the sum of the current and number of rows affected values.

27.7.7.39 mysql_kill()

```
int mysql_kill(MYSQL *mysql, unsigned long pid)
```

Description

Note

`mysql_kill()` is deprecated and will be removed in a future version of MySQL. Instead, use `mysql_query()` to execute a `KILL` statement.

Asks the server to kill the thread specified by `pid`.

This function is deprecated. It is preferable to use `mysql_query()` to issue an SQL `KILL` statement instead.

`mysql_kill()` cannot handle values larger than 32 bits, but to guard against killing the wrong thread returns an error in these cases:

- If given an ID larger than 32 bits, `mysql_kill()` returns a `CR_INVALID_CONN_HANDLE` error.
- After the server's internal thread ID counter reaches a value larger than 32 bits, it returns an `ER_DATA_OUT_OF_RANGE` error for any `mysql_kill()` invocation and `mysql_kill()` fails.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_INVALID_CONN_HANDLE`

  The `pid` was larger than 32 bits.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

- `ER_DATA_OUT_OF_RANGE`

  The server's internal thread ID counter has reached a value larger than 32 bits, at which point it rejects all `mysql_kill()` invocations.

27.7.7.40 mysql_library_end()

```
void mysql_library_end(void)
```

Description

This function finalizes the MySQL library. Call it when you are done using the library (for example, after disconnecting from the server). `mysql_library_end()` performs some memory management to clean up.

For usage information, see Section 27.7.6, “C API Function Overview”, and Section 27.7.7.41, “mysql_library_init()”.

27.7.7.41 mysql_library_init()

```
int mysql_library_init(int argc, char **argv, char **groups)
```

Description

Call this function to initialize the MySQL library before you call any other MySQL function, whether your application is a regular client program or uses the embedded server. If the application uses the embedded server, this call starts the server and initializes any subsystems (`mysys`, `InnoDB`, and so forth) that the server uses.

After your application is done using the MySQL library, call `mysql_library_end()` to clean up. See Section 27.7.7.40, “mysql_library_end()”.

In a nonmulti-threaded environment, the call to `mysql_library_init()` may be omitted, because `mysql_init()` will invoke it automatically as necessary. However, `mysql_library_init()` is not thread-safe in a multi-threaded environment, and thus neither is `mysql_init()`, which calls `mysql_library_init()`. You must either call `mysql_library_init()` prior to spawning any threads, or else use a mutex to protect the call, whether you invoke `mysql_library_init()` or indirectly through `mysql_init()`. Do this prior to any other client library call.

The `argc` and `argv` arguments are analogous to the arguments to `main()`, and enable passing of options to the embedded server. For convenience, `argc` may be `0` (zero) if there are no command-line arguments for the server. This is the usual case for applications intended for use only as regular (nonembedded) clients, and the call typically is written as `mysql_library_init(0, NULL, NULL)`.

```
#include <mysql.h>
#include <stdlib.h>

int main(void) {
  if (mysql_library_init(0, NULL, NULL)) {
    fprintf(stderr, "could not initialize MySQL library\n");
```

```
    exit(1);
  }

  /* Use any MySQL API functions here */

  mysql_library_end();

  return EXIT_SUCCESS;
}
```

When arguments are to be passed (`argc` is greater than `0`), the first element of `argv` is ignored (it typically contains the program name). `mysql_library_init()` makes a copy of the arguments so it is safe to destroy `argv` or `groups` after the call.

For embedded applications, if you want to connect to an external server without starting the embedded server, you have to specify a negative value for `argc`.

The `groups` argument is an array of strings that indicate the groups in option files from which to read options. See Section 4.2.6, “Using Option Files”. Make the final entry in the array `NULL`. For convenience, if the `groups` argument itself is `NULL`, the `[server]` and `[embedded]` groups are used by default.

```
#include <mysql.h>
#include <stdlib.h>

static char *server_args[] = {
  "this_program",       /* this string is not used */
  "--datadir=.",
  "--key_buffer_size=32M"
};
static char *server_groups[] = {
  "embedded",
  "server",
  "this_program_SERVER",
  (char *)NULL
};

int main(void) {
  if (mysql_library_init(sizeof(server_args) / sizeof(char *),
                         server_args, server_groups)) {
    fprintf(stderr, "could not initialize MySQL library\n");
    exit(1);
  }

  /* Use any MySQL API functions here */

  mysql_library_end();

  return EXIT_SUCCESS;
}
```

Return Values

Zero for success. Nonzero if an error occurred.

27.7.7.42 mysql_list_dbs()

```
MYSQL_RES *mysql_list_dbs(MYSQL *mysql, const char *wild)
```

Description

Returns a result set consisting of database names on the server that match the simple regular expression specified by the `wild` parameter. `wild` may contain the wildcard characters `%` or `_`, or may be a `NULL` pointer to match all databases. Calling `mysql_list_dbs()` is similar to executing the query `SHOW DATABASES [LIKE wild]`.

You must free the result set with `mysql_free_result()`.

Return Values

A `MYSQL_RES` result set for success. `NULL` if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.43 mysql_list_fields()

```
MYSQL_RES *mysql_list_fields(MYSQL *mysql, const char *table, const char
*wild)
```

Description

Note

`mysql_list_fields()` is deprecated and will be removed in a future version of MySQL. Instead, use `mysql_query()` to execute a `SHOW COLUMNS` statement.

Returns an empty result set for which the metadata provides information about the columns in the given table that match the simple regular expression specified by the `wild` parameter. `wild` may contain the wildcard characters `%` or `_`, or may be a `NULL` pointer to match all fields. Calling `mysql_list_fields()` is similar to executing the query `SHOW COLUMNS FROM tbl_name [LIKE wild]`.

It is preferable to use `SHOW COLUMNS FROM tbl_name` instead of `mysql_list_fields()`.

You must free the result set with `mysql_free_result()`.

Return Values

A `MYSQL_RES` result set for success. `NULL` if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

```
int i;
MYSQL_RES *tbl_cols = mysql_list_fields(mysql, "mytbl", "f%");

unsigned int field_cnt = mysql_num_fields(tbl_cols);
printf("Number of columns: %d\n", field_cnt);

for (i=0; i < field_cnt; ++i)
{
  /* col describes i-th column of the table */
  MYSQL_FIELD *col = mysql_fetch_field_direct(tbl_cols, i);
  printf ("Column %d: %s\n", i, col->name);
}
mysql_free_result(tbl_cols);
```

27.7.7.44 mysql_list_processes()

```
MYSQL_RES *mysql_list_processes(MYSQL *mysql)
```

Description

Note

`mysql_list_processes()` is deprecated and will be removed in a future version of MySQL. Instead, use `mysql_query()` to execute a `SHOW PROCESSLIST` statement.

Returns a result set describing the current server threads. This is the same kind of information as that reported by `mysqladmin processlist` or a `SHOW PROCESSLIST` query.

You must free the result set with `mysql_free_result()`.

Return Values

A `MYSQL_RES` result set for success. `NULL` if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.45 mysql_list_tables()

```
MYSQL_RES *mysql_list_tables(MYSQL *mysql, const char *wild)
```

Description

Returns a result set consisting of table names in the current database that match the simple regular expression specified by the `wild` parameter. `wild` may contain the wildcard characters `%` or `_`, or may be a `NULL` pointer to match all tables. Calling `mysql_list_tables()` is similar to executing the query `SHOW TABLES [LIKE wild]`.

You must free the result set with `mysql_free_result()`.

Return Values

A `MYSQL_RES` result set for success. `NULL` if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.46 mysql_more_results()

```
bool mysql_more_results(MYSQL *mysql)
```

Description

This function is used when you execute multiple statements specified as a single statement string, or when you execute `CALL` statements, which can return multiple result sets.

`mysql_more_results()` true if more results exist from the currently executed statement, in which case the application must call `mysql_next_result()` to fetch the results.

Return Values

`TRUE` (1) if more results exist. `FALSE` (0) if no more results exist.

In most cases, you can call `mysql_next_result()` instead to test whether more results exist and initiate retrieval if so.

See Section 27.7.20, “C API Support for Multiple Statement Execution”, and Section 27.7.7.47, “mysql_next_result()”.

Errors

None.

27.7.7.47 mysql_next_result()

```
int mysql_next_result(MYSQL *mysql)
```

Description

This function is used when you execute multiple statements specified as a single statement string, or when you use `CALL` statements to execute stored procedures, which can return multiple result sets.

`mysql_next_result()` reads the next statement result and returns a status to indicate whether more results exist. If `mysql_next_result()` returns an error, there are no more results.

Before each call to `mysql_next_result()`, you must call `mysql_free_result()` for the current statement if it is a statement that returned a result set (rather than just a result status).

After calling `mysql_next_result()` the state of the connection is as if you had called `mysql_real_query()` or `mysql_query()` for the next statement. This means that you can call `mysql_store_result()`, `mysql_warning_count()`, `mysql_affected_rows()`, and so forth.

If your program uses `CALL` statements to execute stored procedures, the `CLIENT_MULTI_RESULTS` flag must be enabled. This is because each `CALL` returns a result to indicate the call status, in addition to any result sets that might be returned by statements executed within the procedure. Because `CALL` can return multiple results, process them using a loop that calls `mysql_next_result()` to determine whether there are more results.

`CLIENT_MULTI_RESULTS` can be enabled when you call `mysql_real_connect()`, either explicitly by passing the `CLIENT_MULTI_RESULTS` flag itself, or implicitly by passing `CLIENT_MULTI_STATEMENTS` (which also enables `CLIENT_MULTI_RESULTS`). `CLIENT_MULTI_RESULTS` is enabled by default.

It is also possible to test whether there are more results by calling `mysql_more_results()`. However, this function does not change the connection state, so if it returns true, you must still call `mysql_next_result()` to advance to the next result.

For an example that shows how to use `mysql_next_result()`, see Section 27.7.20, “C API Support for Multiple Statement Execution”.

Return Values

Return Value	Description
0	Successful and there are more results
-1	Successful and there are no more results
>0	An error occurred

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order. For example, if you did not call `mysql_use_result()` for a previous result set.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.48 mysql_num_fields()

```
unsigned int mysql_num_fields(MYSQL_RES *result)
```

To pass a `MYSQL*` argument instead, use `unsigned int mysql_field_count(MYSQL *mysql)`.

Description

Returns the number of columns in a result set.

You can get the number of columns either from a pointer to a result set or to a connection handler. You would use the connection handler if `mysql_store_result()` or `mysql_use_result()` returned `NULL` (and thus you have no result set pointer). In this case, you can call `mysql_field_count()` to determine whether `mysql_store_result()` should have produced a nonempty result. This enables the client program to take proper action without knowing whether the query was a `SELECT` (or `SELECT`-like) statement. The example shown here illustrates how this may be done.

See Section 27.7.18.1, “Why mysql_store_result() Sometimes Returns NULL After mysql_query() Returns Success”.

Return Values

An unsigned integer representing the number of columns in a result set.

Errors

None.

Example

```
MYSQL_RES *result;
unsigned int num_fields;
unsigned int num_rows;

if (mysql_query(&mysql,query_string))
{
    // error
}
else // query succeeded, process any data returned by it
{
    result = mysql_store_result(&mysql);
    if (result)  // there are rows
    {
        num_fields = mysql_num_fields(result);
        // retrieve rows, then call mysql_free_result(result)
    }
    else  // mysql_store_result() returned nothing; should it have?
    {
        if (mysql_errno(&mysql))
        {
           fprintf(stderr, "Error: %s\n", mysql_error(&mysql));
        }
        else if (mysql_field_count(&mysql) == 0)
        {
            // query does not return data
            // (it was not a SELECT)
            num_rows = mysql_affected_rows(&mysql);
        }
    }
}
```

An alternative (if you know that your query should have returned a result set) is to replace the `mysql_errno(&mysql)` call with a check whether `mysql_field_count(&mysql)` returns 0. This happens only if something went wrong.

27.7.7.49 mysql_num_rows()

```
my_ulonglong mysql_num_rows(MYSQL_RES *result)
```

Description

Returns the number of rows in the result set.

The use of `mysql_num_rows()` depends on whether you use `mysql_store_result()` or `mysql_use_result()` to return the result set. If you use `mysql_store_result()`, `mysql_num_rows()` may be called immediately. If you use `mysql_use_result()`, `mysql_num_rows()` does not return the correct value until all the rows in the result set have been retrieved.

`mysql_num_rows()` is intended for use with statements that return a result set, such as `SELECT`. For statements such as `INSERT`, `UPDATE`, or `DELETE`, the number of affected rows can be obtained with `mysql_affected_rows()`.

Return Values

The number of rows in the result set.

Errors

None.

27.7.7.50 mysql_options()

```
int mysql_options(MYSQL *mysql, enum mysql_option option, const void *arg)
```

Description

Can be used to set extra connect options and affect behavior for a connection. This function may be called multiple times to set several options. (To retrieve option values, use `mysql_get_option()`.)

Call `mysql_options()` after `mysql_init()` and before `mysql_connect()` or `mysql_real_connect()`.

The `option` argument is the option that you want to set; the `arg` argument is the value for the option. If the option is an integer, specify a pointer to the value of the integer as the `arg` argument.

The following list describes the possible options, their effect, and how `arg` is used for each option. For option descriptions that indicate `arg` is unused, its value is irrelevant; it is conventional to pass 0.

Note

Several of the options apply only when the application is linked against the `libmysqld` embedded server library and are unused for applications linked against the `libmysqlclient` client library. The embedded server library was removed in MySQL 8.0, so these options are obsolete.

- `MYSQL_DEFAULT_AUTH` (argument type: `char *`)

  The name of the authentication plugin to use.

- `MYSQL_ENABLE_CLEARTEXT_PLUGIN` (argument type: `bool *`)

  Enable the `mysql_clear_password` cleartext authentication plugin. (See Section 6.5.1.3, “Client-Side Cleartext Pluggable Authentication”.)

- `MYSQL_INIT_COMMAND` (argument type: `char *`)

  SQL statement to execute when connecting to the MySQL server. Automatically re-executed if reconnection occurs.

- `MYSQL_OPT_BIND` (argument: `char *`)

  The network interface from which to connect to the server. This is used when the client host has multiple network interfaces. The argument is a host name or IP address (specified as a string).

- `MYSQL_OPT_CAN_HANDLE_EXPIRED_PASSWORDS` (argument type: `bool *`)

  Indicate whether the client can handle expired passwords. For more information, see Section 6.3.9, “Password Expiration and Sandbox Mode”.

- `MYSQL_OPT_COMPRESS` (argument: not used)

  Use the compressed client/server protocol.

- `MYSQL_OPT_CONNECT_ATTR_DELETE` (argument type: `char *`)

  Given a key name, this option deletes a key/value pair from the current set of connection attributes to pass to the server at connect time. The argument is a pointer to a null-terminated string naming the key. Comparison of the key name with existing keys is case sensitive.

  See also the description for the `MYSQL_OPT_CONNECT_ATTR_RESET` option, as well as the description for the `MYSQL_OPT_CONNECT_ATTR_ADD` option in the description of the `mysql_options4()` function. That function description also includes a usage example.

  The Performance Schema exposes connection attributes through the `session_connect_attrs` and `session_account_connect_attrs` tables. See Section 25.11.9, “Performance Schema Connection Attribute Tables”.

- `MYSQL_OPT_CONNECT_ATTR_RESET` (argument not used)

  This option resets (clears) the current set of connection attributes to pass to the server at connect time.

  See also the description for the `MYSQL_OPT_CONNECT_ATTR_DELETE` option, as well as the description for the `MYSQL_OPT_CONNECT_ATTR_ADD` option in the description of the `mysql_options4()` function. That function description also includes a usage example.

  The Performance Schema exposes connection attributes through the `session_connect_attrs` and `session_account_connect_attrs` tables. See Section 25.11.9, “Performance Schema Connection Attribute Tables”.

- `MYSQL_OPT_CONNECT_TIMEOUT` (argument type: `unsigned int *`)

  The connect timeout in seconds.

- `MYSQL_OPT_GUESS_CONNECTION` (argument: not used)

  Note

  The embedded server library was removed in MySQL 8.0, so this option is obsolete.

  For an application linked against the `libmysqld` embedded server library, this enables the library to guess whether to use the embedded server or a remote server. “Guess” means that if the host name is set and is not `localhost`, it uses a remote server. This behavior is the default. `MYSQL_OPT_USE_EMBEDDED_CONNECTION` and `MYSQL_OPT_USE_REMOTE_CONNECTION` can be used to override it. This option is ignored for applications linked against the `libmysqlclient` client library.

- `MYSQL_OPT_LOCAL_INFILE` (argument type: optional pointer to `unsigned int`)

  This option affects client-side `LOCAL` capability for `LOAD DATA` operations. By default, `LOCAL` capability is determined by the default compiled into the MySQL client library (see Section 13.2.7, “LOAD DATA INFILE Syntax”). To control this capability explicitly, invoke `mysql_options()` to set the `MYSQL_OPT_LOCAL_INFILE` option:

  - `LOCAL` is disabled if the pointer points to an `unsigned int` that has a zero value.

  - `LOCAL` is enabled if no pointer is given or if the pointer points to an `unsigned int` that has a nonzero value.

  Successful use of a `LOCAL` load operation by a client also requires that the server permits it.

- `MYSQL_OPT_MAX_ALLOWED_PACKET` (argument: `unsigned long *`)

  This option sets the `max_allowed_packet` system variable. If the `mysql` argument is non-`NULL`, the call sets the session system variable value for that session. If `mysql` is `NULL`, the call sets the global system variable value.

- `MYSQL_OPT_NAMED_PIPE` (argument: not used)

  Use a named pipe to connect to the MySQL server on Windows, if the server permits named-pipe connections.

- `MYSQL_OPT_NET_BUFFER_LENGTH` (argument: `unsigned long *`)

  This option sets the `net_buffer_length` system variable. If the `mysql` argument is non-`NULL`, the call sets the session system variable value for that session. If `mysql` is `NULL`, the call sets the global system variable value.

- `MYSQL_OPT_PROTOCOL` (argument type: `unsigned int *`)

  Type of protocol to use. Specify one of the enum values of `mysql_protocol_type` defined in `mysql.h`.

- `MYSQL_OPT_READ_TIMEOUT` (argument type: `unsigned int *`)

  The timeout in seconds for each attempt to read from the server. There are retries if necessary, so the total effective timeout value is three times the option value. You can set the value so that a lost connection can be detected earlier than the TCP/IP `Close_Wait_Timeout` value of 10 minutes.

- `MYSQL_OPT_RECONNECT` (argument type: `bool *`)

  Enable or disable automatic reconnection to the server if the connection is found to have been lost. Reconnect is off by default; this option provides a way to set reconnection behavior explicitly.

- `MYSQL_OPT_RETRY_COUNT` (argument type: `unsigned int *`)

  The retry count for I/O-related system calls that are interrupted while connecting to the server or communicating with it. The default value is 1 (1 retry if the initial call is interrupted for 2 tries total).

- `MYSQL_OPT_SSL_CA` (argument type: `char *`)

  The path to a file in PEM format that contains a list of trusted SSL CAs.

- `MYSQL_OPT_SSL_CAPATH` (argument type: `char *`)

  The path to a directory that contains trusted SSL CA certificates in PEM format.

- `MYSQL_OPT_SSL_CERT` (argument type: `char *`)

  The name of an SSL certificate file in PEM format to use for establishing a secure connection.

- `MYSQL_OPT_SSL_CIPHER` (argument type: `char *`)

  A list of permissible ciphers to use for SSL encryption.

- `MYSQL_OPT_SSL_CRL` (argument type: `char *`)

  The path to a file containing certificate revocation lists in PEM format.

- `MYSQL_OPT_SSL_CRLPATH` (argument type: `char *`)

  The path to a directory that contains files containing certificate revocation lists in PEM format.

- `MYSQL_OPT_SSL_KEY` (argument type: `char *`)

  The name of an SSL key file in PEM format to use for establishing a secure connection.

- `MYSQL_OPT_SSL_MODE` (argument type: `unsigned int *`)

  The security state to use for the connection to the server: `SSL_MODE_DISABLED`, `SSL_MODE_PREFERRED`, `SSL_MODE_REQUIRED`, `SSL_MODE_VERIFY_CA`, `SSL_MODE_VERIFY_IDENTITY`. The default is `SSL_MODE_PREFERRED`. These modes are the permitted values of the `mysql_ssl_mode` enumeration defined in `mysql.h`. For more information about the security states, see the description of `--ssl-mode` in Section 6.4.5, “Command Options for Secure Connections”.

- `MYSQL_OPT_TLS_VERSION` (argument type: `char *`)

  The protocols permitted by the client for encrypted connections. The value is a comma-separated list containing one or more protocol names. The protocols that can be named for this option depend on the SSL library used to compile MySQL. For details, see Section 6.4.3, “Secure Connection Protocols and Ciphers”.

- `MYSQL_OPT_USE_EMBEDDED_CONNECTION` (argument: not used)

  Note

  The embedded server library was removed in MySQL 8.0, so this option is obsolete.

  For an application linked against the `libmysqld` embedded server library, this forces the use of the embedded server for the connection. This option is ignored for applications linked against the `libmysqlclient` client library.

- `MYSQL_OPT_USE_REMOTE_CONNECTION` (argument: not used)

  Note

  The embedded server library was removed in MySQL 8.0, so this option is obsolete.

  For an application linked against the `libmysqld` embedded server library, this forces the use of a remote server for the connection. This option is ignored for applications linked against the `libmysqlclient` client library.

- `MYSQL_OPT_USE_RESULT` (argument: not used)

  This option is unused.

- `MYSQL_OPT_WRITE_TIMEOUT` (argument type: `unsigned int *`)

  The timeout in seconds for each attempt to write to the server. There is a retry if necessary, so the total effective timeout value is two times the option value.

- `MYSQL_PLUGIN_DIR` (argument type: `char *`)

  The directory in which to look for client plugins.

- `MYSQL_READ_DEFAULT_FILE` (argument type: `char *`)

  Read options from the named option file instead of from `my.cnf`.

- `MYSQL_READ_DEFAULT_GROUP` (argument type: `char *`)

  Read options from the named group from `my.cnf` or the file specified with `MYSQL_READ_DEFAULT_FILE`.

- `MYSQL_REPORT_DATA_TRUNCATION` (argument type: `bool *`)

  Enable or disable reporting of data truncation errors for prepared statements using the `error` member of `MYSQL_BIND` structures. (Default: enabled.)

- `MYSQL_SECURE_AUTH` (argument type: `bool *`)

  Whether to connect to a server that does not support the password hashing used in MySQL 4.1.1 and later. This option is enabled by default.

- `MYSQL_SERVER_PUBLIC_KEY` (argument type: `char *`)

  The path name to a file containing the server RSA public key. The file must be in PEM format. The public key is used for RSA encryption of the client password for connections to the server made using accounts that authenticate with the `sha256_password` plugin. This option is ignored for client accounts that do not authenticate with that plugin. It is also ignored if password encryption is not needed, as is the case when the client connects to the server using an SSL connection.

  The server sends the public key to the client as needed, so it is not necessary to use this option for RSA password encryption to occur. It is more efficient to do so because then the server need not send the key.

  For additional discussion regarding use of the `sha256_password` plugin, including how to get the RSA public key, see Section 6.5.1.2, “SHA-256 Pluggable Authentication”.

- `MYSQL_SET_CHARSET_DIR` (argument type: `char *`)

  The path name to the directory that contains character set definition files.

- `MYSQL_SET_CHARSET_NAME` (argument type: `char *`)

  The name of the character set to use as the default character set. The argument can be `MYSQL_AUTODETECT_CHARSET_NAME` to cause the character set to be autodetected based on the operating system setting (see Section 10.1.4, “Connection Character Sets and Collations”).

- `MYSQL_SET_CLIENT_IP` (argument type: `char *`)

  Note

  The embedded server library was removed in MySQL 8.0, so this option is obsolete.

  For an application linked against the `libmysqld` embedded server library (when `libmysqld` is compiled with authentication support), this means that the user is considered to have connected from the specified IP address (specified as a string) for authentication purposes. This option is ignored for applications linked against the `libmysqlclient` client library.

- `MYSQL_SHARED_MEMORY_BASE_NAME` (argument type: `char *`)

  The name of the shared-memory object for communication to the server on Windows, if the server supports shared-memory connections. Specify the same value as the `--shared-memory-base-name` option used for the `mysqld` server you want to connect to.

The `client` group is always read if you use `MYSQL_READ_DEFAULT_FILE` or `MYSQL_READ_DEFAULT_GROUP`.

The specified group in the option file may contain the following options.

Option	Description
`character-sets-dir=dir_name`	The directory where character sets are installed.
`compress`	Use the compressed client/server protocol.
`connect-timeout=seconds`	The connect timeout in seconds. On Linux this timeout is also used for waiting for the first answer from the server.
`database=db_name`	Connect to this database if no database was specified in the connect command.
`debug`	Debug options.
`default-character-set=charset_name`	The default character set to use.
`disable-local-infile`	Disable use of `LOAD DATA LOCAL INFILE`.
`enable-cleartext-plugin`	Enable the `mysql_clear_password` cleartext authentication plugin.
`host=host_name`	Default host name.
`init-command=stmt`	Statement to execute when connecting to MySQL server. Automatically re-executed if reconnection occurs.
`interactive-timeout=seconds`	Same as specifying `CLIENT_INTERACTIVE` to `mysql_real_connect()`. See Section 27.7.7.54, “mysql_real_connect()”.
`local-infile[={0|1}]`	If no argument or nonzero argument, enable use of `LOAD DATA LOCAL`; otherwise disable.
`max_allowed_packet=bytes`	Maximum size of packet that client can read from server.
`multi-queries`, `multi-results`	Enable multiple result sets from multiple-statement executions or stored procedures.
`multi-statements`	Enable the client to send multiple statements in a single string (separated by `;` characters).
`password=password`	Default password.
`pipe`	Use named pipes to connect to a MySQL server on Windows.
`port=port_num`	Default port number.
`protocol={TCP|SOCKET|PIPE|MEMORY}`	The protocol to use when connecting to the server.
`return-found-rows`	Tell `mysql_info()` to return found rows instead of updated rows when using `UPDATE`.
`shared-memory-base-name=name`	Shared-memory name to use to connect to server.
`socket={file_name|pipe_name}`	Default socket file.
`ssl-ca=file_name`	Certificate Authority file.
`ssl-capath=dir_name`	Certificate Authority directory.
`ssl-cert=file_name`	Certificate file.
`ssl-cipher=cipher_list`	Permissible SSL ciphers.
`ssl-key=file_name`	Key file.
`timeout=seconds`	Like `connect-timeout`.
`user`	Default user.

`timeout` has been replaced by `connect-timeout`, but `timeout` is still supported for backward compatibility.

For more information about option files used by MySQL programs, see Section 4.2.6, “Using Option Files”.

Return Values

Zero for success. Nonzero if you specify an unknown option.

Example

The following `mysql_options()` calls request the use of compression in the client/server protocol, cause options to be read from the `[odbc]` group of option files, and disable transaction autocommit mode:

```
MYSQL mysql;

mysql_init(&mysql);
mysql_options(&mysql,MYSQL_OPT_COMPRESS,0);
mysql_options(&mysql,MYSQL_READ_DEFAULT_GROUP,"odbc");
mysql_options(&mysql,MYSQL_INIT_COMMAND,"SET autocommit=0");
if (!mysql_real_connect(&mysql,"host","user","passwd","database",0,NULL,0))
{
    fprintf(stderr, "Failed to connect to database: Error: %s\n",
          mysql_error(&mysql));
}
```

This code requests that the client use the compressed client/server protocol and read the additional options from the `odbc` section in the `my.cnf` file.

27.7.7.51 mysql_options4()

```
int mysql_options4(MYSQL *mysql, enum mysql_option option, const void *arg1,
const void *arg2)
```

Description

`mysql_options4()` is similar to `mysql_options()` but has an extra fourth argument so that two values can be passed for the option specified in the second argument.

The following list describes the permitted options, their effect, and how `arg1` and `arg2` are used.

- `MYSQL_OPT_CONNECT_ATTR_ADD` (argument types: `char *, char *`)

  This option adds an attribute key/value pair to the current set of connection attributes to pass to the server at connect time. Both arguments are pointers to null-terminated strings. The first and second strings indicate the key and value, respectively. If the key is empty or already exists in the current set of connection attributes, an error occurs. Comparison of the key name with existing keys is case sensitive.

  Key names that begin with an underscore (_) are reserved for internal use and should not be created by application programs. This convention permits new attributes to be introduced by MySQL without colliding with application attributes.

  `mysql_options4()` imposes a limit of 64KB on the aggregate size of connection attribute data it will accept. For calls that cause this limit to be exceeded, a `CR_INVALID_PARAMETER_NO` error occurs. Attribute size-limit checks also occur on the server side. For details, see Section 25.11.9, “Performance Schema Connection Attribute Tables”, which also describes how the Performance Schema exposes connection attributes through the `session_connect_attrs` and `session_account_connect_attrs` tables.

  See also the descriptions for the `MYSQL_OPT_CONNECT_ATTR_RESET` `MYSQL_OPT_CONNECT_ATTR_DELETE` options in the description of the `mysql_options()` function.

Return Values

Zero for success. Nonzero if you specify an unknown option.

Errors

- `CR_DUPLICATE_CONNECTION_ATTR`

  A duplicate attribute name was specified.

- `CR_INVALID_PARAMETER_NO`

  A key name was empty or the amount of key/value connection attribute data exceeds 64KB limit.

- `CR_OUT_OF_MEMORY`

  Out of memory.

Example

This example demonstrates the calls that specify connection attributes:

```
MYSQL mysql;

mysql_init(&mysql);
mysql_options(&mysql,MYSQL_OPT_CONNECT_ATTR_RESET, 0);
mysql_options4(&mysql,MYSQL_OPT_CONNECT_ATTR_ADD, "key1", "value1");
mysql_options4(&mysql,MYSQL_OPT_CONNECT_ATTR_ADD, "key2", "value2");
mysql_options4(&mysql,MYSQL_OPT_CONNECT_ATTR_ADD, "key3", "value3");
mysql_options(&mysql,MYSQL_OPT_CONNECT_ATTR_DELETE, "key1");
if (!mysql_real_connect(&mysql,"host","user","passwd","database",0,NULL,0))
{
    fprintf(stderr, "Failed to connect to database: Error: %s\n",
        mysql_error(&mysql));
}
```

27.7.7.52 mysql_ping()

```
int mysql_ping(MYSQL *mysql)
```

Description

Checks whether the connection to the server is working. If the connection has gone down and auto-reconnect is enabled an attempt to reconnect is made. If the connection is down and auto-reconnect is disabled, `mysql_ping()` returns an error.

Auto-reconnect is disabled by default. To enable it, call `mysql_options()` with the `MYSQL_OPT_RECONNECT` option. For details, see Section 27.7.7.50, “mysql_options()”.

`mysql_ping()` can be used by clients that remain idle for a long while, to check whether the server has closed the connection and reconnect if necessary.

If `mysql_ping()`) does cause a reconnect, there is no explicit indication of it. To determine whether a reconnect occurs, call `mysql_thread_id()` to get the original connection identifier before calling `mysql_ping()`, then call `mysql_thread_id()` again to see whether the identifier has changed.

If reconnect occurs, some characteristics of the connection will have been reset. For details about these characteristics, see Section 27.7.19, “Controlling Automatic Reconnection Behavior”.

Return Values

Zero if the connection to the server is active. Nonzero if an error occurred. A nonzero return does not indicate whether the MySQL server itself is down; the connection might be broken for other reasons such as network problems.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.53 mysql_query()

```
int mysql_query(MYSQL *mysql, const char *stmt_str)
```

Description

Executes the SQL statement pointed to by the null-terminated string `stmt_str`. Normally, the string must consist of a single SQL statement without a terminating semicolon (`;`) or `\g`. If multiple-statement execution has been enabled, the string can contain several statements separated by semicolons. See Section 27.7.20, “C API Support for Multiple Statement Execution”.

`mysql_query()` cannot be used for statements that contain binary data; you must use `mysql_real_query()` instead. (Binary data may contain the `\0` character, which `mysql_query()` interprets as the end of the statement string.)

If you want to know whether the statement returns a result set, you can use `mysql_field_count()` to check for this. See Section 27.7.7.22, “mysql_field_count()”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.54 mysql_real_connect()

```
MYSQL *mysql_real_connect(MYSQL *mysql, const char *host, const char
*user, const char *passwd, const char *db, unsigned int port, const char
*unix_socket, unsigned long client_flag)
```

Description

`mysql_real_connect()` attempts to establish a connection to a MySQL database engine running on `host`. `mysql_real_connect()` must complete successfully before you can execute any other API functions that require a valid `MYSQL` connection handler structure.

The parameters are specified as follows:

- For the first parameter, specify the address of an existing `MYSQL` structure. Before calling `mysql_real_connect()`, call `mysql_init()` to initialize the `MYSQL` structure. You can change a lot of connect options with the `mysql_options()` call. See Section 27.7.7.50, “mysql_options()”.
- The value of `host` may be either a host name or an IP address. If `host` is `NULL` or the string `"localhost"`, a connection to the local host is assumed. For Windows, the client connects using a shared-memory connection, if the server has shared-memory connections enabled. Otherwise, TCP/IP is used. For Unix, the client connects using a Unix socket file. For local connections, you can also influence the type of connection to use with the `MYSQL_OPT_PROTOCOL` or `MYSQL_OPT_NAMED_PIPE` options to `mysql_options()`. The type of connection must be supported by the server. For a `host` value of `"."` on Windows, the client connects using a named pipe, if the server has named-pipe connections enabled. If named-pipe connections are not enabled, an error occurs.
- The `user` parameter contains the user's MySQL login ID. If `user` is `NULL` or the empty string `""`, the current user is assumed. Under Unix, this is the current login name. Under Windows ODBC, the current user name must be specified explicitly. See the Connector/ODBC section of Chapter 27, *Connectors and APIs*.
- The `passwd` parameter contains the password for `user`. If `passwd` is `NULL`, only entries in the `user` table for the user that have a blank (empty) password field are checked for a match. This enables the database administrator to set up the MySQL privilege system in such a way that users get different privileges depending on whether they have specified a password.

  Note

  Do not attempt to encrypt the password before calling `mysql_real_connect()`; password encryption is handled automatically by the client API.
- The `user` and `passwd` parameters use whatever character set has been configured for the `MYSQL` object. By default, this is `utf8mb4` (`latin1` prior to MySQL 8.0.2), but can be changed by calling `mysql_options(mysql, MYSQL_SET_CHARSET_NAME, "charset_name")` prior to connecting.
- `db` is the database name. If `db` is not `NULL`, the connection sets the default database to this value.
- If `port` is not 0, the value is used as the port number for the TCP/IP connection. Note that the `host` parameter determines the type of the connection.
- If `unix_socket` is not `NULL`, the string specifies the socket or named pipe to use. Note that the `host` parameter determines the type of the connection.
- The value of `client_flag` is usually 0, but can be set to a combination of the following flags to enable certain features.

Flag Name	Flag Description
`CAN_HANDLE_EXPIRED_PASSWORDS`	The client can handle expired passwords. For more information, see Section 6.3.9, “Password Expiration and Sandbox Mode”.
`CLIENT_COMPRESS`	Use compression in the client/server protocol.
`CLIENT_FOUND_ROWS`	Return the number of found (matched) rows, not the number of changed rows.
`CLIENT_IGNORE_SIGPIPE`	Prevents the client library from installing a `SIGPIPE` signal handler. This can be used to avoid conflicts with a handler that the application has already installed.
`CLIENT_IGNORE_SPACE`	Permit spaces after function names. Makes all functions names reserved words.

Flag Name	Flag Description
CLIENT_INTERACTIVE	Permit interactive_timeout seconds of inactivity (rather than wait_timeout seconds) before closing the connection. The client's session wait_timeout variable is set to the value of the session interactive_timeout variable.
CLIENT_LOCAL_FILES	Enable LOAD DATA LOCAL handling.
CLIENT_MULTI_RESULTS	Tell the server that the client can handle multiple result sets from multiple-statement executions or stored procedures. This flag is automatically enabled if CLIENT_MULTI_STATEMENTS is enabled. See the note following this table for more information about this flag.
CLIENT_MULTI_STATEMENTS	Tell the server that the client may send multiple statements in a single string (separated by ; characters). If this flag is not set, multiple-statement execution is disabled. See the note following this table for more information about this flag.
CLIENT_NO_SCHEMA	Do not permit *db_name.tbl_name.col_name* syntax. This is for ODBC. It causes the parser to generate an error if you use that syntax, which is useful for trapping bugs in some ODBC programs.
CLIENT_ODBC	Unused.
CLIENT_SSL	Use SSL (encrypted protocol). Do not set this option within an application program; it is set internally in the client library. Instead, use mysql_ssl_set() before calling mysql_real_connect().
CLIENT_REMEMBER_OPTIONS	Remember options specified by calls to mysql_options(). Without this option, if mysql_real_connect() fails, you must repeat the mysql_options() calls before trying to connect again. With this option, the mysql_options() calls need not be repeated.

If your program uses CALL statements to execute stored procedures, the CLIENT_MULTI_RESULTS flag must be enabled. This is because each CALL returns a result to indicate the call status, in addition to any result sets that might be returned by statements executed within the procedure. Because CALL can return multiple results, process them using a loop that calls mysql_next_result() to determine whether there are more results.

CLIENT_MULTI_RESULTS can be enabled when you call mysql_real_connect(), either explicitly by passing the CLIENT_MULTI_RESULTS flag itself, or implicitly by passing CLIENT_MULTI_STATEMENTS (which also enables CLIENT_MULTI_RESULTS). CLIENT_MULTI_RESULTS is enabled by default.

If you enable CLIENT_MULTI_STATEMENTS or CLIENT_MULTI_RESULTS, process the result for every call to mysql_query() or mysql_real_query() by using a loop that calls mysql_next_result() to determine whether there are more results. For an example, see Section 27.7.20, “C API Support for Multiple Statement Execution”.

For some parameters, it is possible to have the value taken from an option file rather than from an explicit value in the mysql_real_connect() call. To do this, call mysql_options() with the MYSQL_READ_DEFAULT_FILE or MYSQL_READ_DEFAULT_GROUP option before calling mysql_real_connect(). Then, in the mysql_real_connect() call, specify the “no-value” value for each parameter to be read from an option file:

- For host, specify a value of NULL or the empty string ("").
- For user, specify a value of NULL or the empty string.

- For `passwd`, specify a value of `NULL`. (For the password, a value of the empty string in the `mysql_real_connect()` call cannot be overridden in an option file, because the empty string indicates explicitly that the MySQL account must have an empty password.)
- For `db`, specify a value of `NULL` or the empty string.
- For `port`, specify a value of 0.
- For `unix_socket`, specify a value of `NULL`.

If no value is found in an option file for a parameter, its default value is used as indicated in the descriptions given earlier in this section.

Return Values

A `MYSQL*` connection handler if the connection was successful, `NULL` if the connection was unsuccessful. For a successful connection, the return value is the same as the value of the first parameter.

Errors

- `CR_CONN_HOST_ERROR`

  Failed to connect to the MySQL server.

- `CR_CONNECTION_ERROR`

  Failed to connect to the local MySQL server.

- `CR_IPSOCK_ERROR`

  Failed to create an IP socket.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SOCKET_CREATE_ERROR`

  Failed to create a Unix socket.

- `CR_UNKNOWN_HOST`

  Failed to find the IP address for the host name.

- `CR_VERSION_ERROR`

  A protocol mismatch resulted from attempting to connect to a server with a client library that uses a different protocol version.

- `CR_NAMEDPIPEOPEN_ERROR`

  Failed to create a named pipe on Windows.

- `CR_NAMEDPIPEWAIT_ERROR`

  Failed to wait for a named pipe on Windows.

- `CR_NAMEDPIPESETSTATE_ERROR`

  Failed to get a pipe handler on Windows.

- `CR_SERVER_LOST`

If `connect_timeout` > 0 and it took longer than `connect_timeout` seconds to connect to the server or if the server died while executing the `init-command`.

- `CR_ALREADY_CONNECTED`

  The `MYSQL` connection handler is already connected.

Example

```
MYSQL mysql;

mysql_init(&mysql);
mysql_options(&mysql,MYSQL_READ_DEFAULT_GROUP,"your_prog_name");
if (!mysql_real_connect(&mysql,"host","user","passwd","database",0,NULL,0))
{
    fprintf(stderr, "Failed to connect to database: Error: %s\n",
          mysql_error(&mysql));
}
```

By using `mysql_options()` the MySQL library reads the `[client]` and `[your_prog_name]` sections in the `my.cnf` file which ensures that your program works, even if someone has set up MySQL in some nonstandard way.

Upon connection, `mysql_real_connect()` sets the `reconnect` flag (part of the `MYSQL` structure) to a value of `1` in versions of the API older than 5.0.3, or `0` in newer versions. A value of `1` for this flag indicates that if a statement cannot be performed because of a lost connection, to try reconnecting to the server before giving up. You can use the `MYSQL_OPT_RECONNECT` option to `mysql_options()` to control reconnection behavior.

27.7.7.55 mysql_real_escape_string()

```
unsigned long mysql_real_escape_string(MYSQL *mysql, char *to, const char
*from, unsigned long length)
```

Description

This function creates a legal SQL string for use in an SQL statement. See Section 9.1.1, “String Literals”.

Note

`mysql_real_escape_string()` fails and produces an `CR_INSECURE_API_ERR` error if the `NO_BACKSLASH_ESCAPES` SQL mode is enabled. In this case, the function cannot escape quote characters except by doubling them, and to do this properly, it must know more information about the quoting context than is available. Instead, use `mysql_real_escape_string_quote()`, which takes an extra argument for specifying the quoting context.

The `mysql` argument must be a valid, open connection because character escaping depends on the character set in use by the server.

The string in the `from` argument is encoded to produce an escaped SQL string, taking into account the current character set of the connection. The result is placed in the `to` argument, followed by a terminating null byte.

Characters encoded are `\`, `'`, `"`, `NUL` (ASCII 0), `\n`, `\r`, and Control+Z. Strictly speaking, MySQL requires only that backslash and the quote character used to quote the string in the query be escaped. `mysql_real_escape_string()` quotes the other characters to make them easier to read in log files. For comparison, see the quoting rules for literal strings and the `QUOTE()` SQL function in Section 9.1.1, “String Literals”, and Section 12.5, “String Functions”.

The string pointed to by `from` must be `length` bytes long. You must allocate the `to` buffer to be at least `length*2+1` bytes long. (In the worst case, each character may need to be encoded as using two bytes, and there must be room for the terminating null byte.) When `mysql_real_escape_string()` returns, the contents of `to` is a null-terminated string. The return value is the length of the encoded string, not including the terminating null byte.

If you must change the character set of the connection, use the `mysql_set_character_set()` function rather than executing a `SET NAMES` (or `SET CHARACTER SET`) statement. `mysql_set_character_set()` works like `SET NAMES` but also affects the character set used by `mysql_real_escape_string()`, which `SET NAMES` does not.

Example

The following example inserts two escaped strings into an `INSERT` statement, each within single quote characters:

```
char query[1000],*end;

end = my_stpcpy(query,"INSERT INTO test_table VALUES('");
end += mysql_real_escape_string(&mysql,end,"What is this",12);
end = my_stpcpy(end,"','");
end += mysql_real_escape_string(&mysql,end,"binary data: \0\r\n",16);
end = my_stpcpy(end,"')");

if (mysql_real_query(&mysql,query,(unsigned int) (end - query)))
{
   fprintf(stderr, "Failed to insert row, Error: %s\n",
           mysql_error(&mysql));
}
```

The `my_stpcpy()` function used in the example is included in the `libmysqlclient` library and works like `strcpy()` but returns a pointer to the terminating null of the first parameter.

Return Values

The length of the encoded string that is placed into the `to` argument, not including the terminating null byte, or -1 if an error occurs.

Because `mysql_real_escape_string()` returns an unsigned value, you can check for -1 by comparing the return value to `(unsigned long)-1` (or to `(unsigned long)~0`, which is equivalent).

Errors

- `CR_INSECURE_API_ERR`

  This error occurs if the `NO_BACKSLASH_ESCAPES` SQL mode is enabled because, in that case, `mysql_real_escape_string()` cannot be guaranteed to produce a properly encoded result. To avoid this error, use `mysql_real_escape_string_quote()` instead.

27.7.7.56 mysql_real_escape_string_quote()

```
unsigned long mysql_real_escape_string_quote(MYSQL *mysql, char *to, const
char *from, unsigned long length, char quote)
```

Description

This function creates a legal SQL string for use in an SQL statement. See Section 9.1.1, “String Literals”.

The `mysql` argument must be a valid, open connection because character escaping depends on the character set in use by the server.

The string in the `from` argument is encoded to produce an escaped SQL string, taking into account the current character set of the connection. The result is placed in the `to` argument, followed by a terminating null byte.

Characters encoded are `\`, `'`, `"`, `NUL` (ASCII 0), `\n`, `\r`, Control+Z, and `` ` ``. Strictly speaking, MySQL requires only that backslash and the quote character used to quote the string in the query be escaped. `mysql_real_escape_string_quote()` quotes the other characters to make them easier to read in log files. For comparison, see the quoting rules for literal strings and the `QUOTE()` SQL function in Section 9.1.1, “String Literals”, and Section 12.5, “String Functions”.

Note

If the `ANSI_QUOTES` SQL mode is enabled, `mysql_real_escape_string_quote()` cannot be used to escape double quote characters for use within double-quoted identifiers. (The function cannot tell whether the mode is enabled to determine the proper escaping character.)

The string pointed to by `from` must be `length` bytes long. You must allocate the `to` buffer to be at least `length*2+1` bytes long. (In the worst case, each character may need to be encoded as using two bytes, and there must be room for the terminating null byte.) When `mysql_real_escape_string_quote()` returns, the contents of `to` is a null-terminated string. The return value is the length of the encoded string, not including the terminating null byte.

The `quote` argument indicates the context in which the escaped string is to be placed. Suppose that you intend to escape the `from` argument and insert the escaped string (designated here by *`str`*) into one of the following statements:

```
1) SELECT * FROM table WHERE name = 'str'
2) SELECT * FROM table WHERE name = "str"
3) SELECT * FROM `str` WHERE id = 103
```

To perform escaping properly for each statement, call `mysql_real_escape_string_quote()` as follows, where the final argument indicates the quoting context:

```
1) len = mysql_real_escape_string_quote(&mysql,to,from,from_len,'\'');
2) len = mysql_real_escape_string_quote(&mysql,to,from,from_len,'"');
3) len = mysql_real_escape_string_quote(&mysql,to,from,from_len,'`');
```

If you must change the character set of the connection, use the `mysql_set_character_set()` function rather than executing a `SET NAMES` (or `SET CHARACTER SET`) statement. `mysql_set_character_set()` works like `SET NAMES` but also affects the character set used by `mysql_real_escape_string_quote()`, which `SET NAMES` does not.

Example

The following example inserts two escaped strings into an `INSERT` statement, each within single quote characters:

```
char query[1000],*end;

end = my_stpcpy(query,"INSERT INTO test_table VALUES('");
end += mysql_real_escape_string_quote(&mysql,end,"What is this",12,'\'');
end = my_stpcpy(end,"','");
end += mysql_real_escape_string_quote(&mysql,end,"binary data: \0\r\n",16,'\'');
end = my_stpcpy(end,"')");

if (mysql_real_query(&mysql,query,(unsigned int) (end - query)))
{
   fprintf(stderr, "Failed to insert row, Error: %s\n",
           mysql_error(&mysql));
}
```

The `my_stpcpy()` function used in the example is included in the `libmysqlclient` library and works like `strcpy()` but returns a pointer to the terminating null of the first parameter.

Return Values

The length of the encoded string that is placed into the `to` argument, not including the terminating null byte.

Errors

None.

27.7.7.57 mysql_real_query()

```
int mysql_real_query(MYSQL *mysql, const char *stmt_str, unsigned long
length)
```

Description

Executes the SQL statement pointed to by `stmt_str`, a string `length` bytes long. Normally, the string must consist of a single SQL statement without a terminating semicolon (`;`) or `\g`. If multiple-statement execution has been enabled, the string can contain several statements separated by semicolons. See Section 27.7.20, “C API Support for Multiple Statement Execution”.

`mysql_query()` cannot be used for statements that contain binary data; you must use `mysql_real_query()` instead. (Binary data may contain the `\0` character, which `mysql_query()` interprets as the end of the statement string.) In addition, `mysql_real_query()` is faster than `mysql_query()` because it does not call `strlen()` on the statement string.

If you want to know whether the statement returns a result set, you can use `mysql_field_count()` to check for this. See Section 27.7.7.22, “mysql_field_count()”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.58 mysql_refresh()

```
int mysql_refresh(MYSQL *mysql, unsigned int options)
```

Description

Note

`mysql_refresh()` is deprecated and will be removed in a future version of MySQL. Instead, use `mysql_query()` to execute a `FLUSH` statement.

This function flushes tables or caches, or resets replication server information. The connected user must have the `RELOAD` privilege.

The `options` argument is a bitmask composed from any combination of the following values. Multiple values can be OR'ed together to perform multiple operations with a single call.

- `REFRESH_GRANT`

  Refresh the grant tables, like `FLUSH PRIVILEGES`.

- `REFRESH_LOG`

  Flush the logs, like `FLUSH LOGS`.

- `REFRESH_TABLES`

  Flush the table cache, like `FLUSH TABLES`.

- `REFRESH_HOSTS`

  Flush the host cache, like `FLUSH HOSTS`.

- `REFRESH_STATUS`

  Reset status variables, like `FLUSH STATUS`.

- `REFRESH_THREADS`

  Flush the thread cache.

- `REFRESH_SLAVE`

  On a slave replication server, reset the master server information and restart the slave, like `RESET SLAVE`.

- `REFRESH_MASTER`

  On a master replication server, remove the binary log files listed in the binary log index and truncate the index file, like `RESET MASTER`.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.59 mysql_reload()

```
int mysql_reload(MYSQL *mysql)
```

Description

Asks the MySQL server to reload the grant tables. The connected user must have the `RELOAD` privilege.

This function is deprecated. It is preferable to use `mysql_query()` to issue an SQL `FLUSH PRIVILEGES` statement instead.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.60 mysql_reset_connection()

```
int mysql_reset_connection(MYSQL *mysql)
```

Description

Resets the connection to clear the session state.

`mysql_reset_connection()` has effects similar to `mysql_change_user()` or an auto-reconnect except that the connection is not closed and reopened, and reauthentication is not done. See Section 27.7.7.3, “mysql_change_user()”) and see Section 27.7.19, “Controlling Automatic Reconnection Behavior”).

The connection-related state is affected as follows:

- Any active transactions are rolled back and autocommit mode is reset.
- All table locks are released.
- All `TEMPORARY` tables are closed (and dropped).
- Session system variables are reinitialized to the values of the corresponding global system variables, including system variables that are set implicitly by statements such as `SET NAMES`.
- User variable settings are lost.
- Prepared statements are released.
- `HANDLER` variables are closed.
- The value of `LAST_INSERT_ID()` is reset to 0.

- Locks acquired with `GET_LOCK()` are released.

Return Values

Zero for success. Nonzero if an error occurred.

27.7.7.61 mysql_rollback()

```
bool mysql_rollback(MYSQL *mysql)
```

Description

Rolls back the current transaction.

The action of this function is subject to the value of the `completion_type` system variable. In particular, if the value of `completion_type` is `RELEASE` (or 2), the server performs a release after terminating a transaction and closes the client connection. Call `mysql_close()` from the client program to close the connection from the client side.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

None.

27.7.7.62 mysql_row_seek()

```
MYSQL_ROW_OFFSET mysql_row_seek(MYSQL_RES *result, MYSQL_ROW_OFFSET offset)
```

Description

Sets the row cursor to an arbitrary row in a query result set. The `offset` value is a row offset, typically a value returned from `mysql_row_tell()` or from `mysql_row_seek()`. This value is not a row number; to seek to a row within a result set by number, use `mysql_data_seek()` instead.

This function requires that the result set structure contains the entire result of the query, so `mysql_row_seek()` may be used only in conjunction with `mysql_store_result()`, not with `mysql_use_result()`.

Return Values

The previous value of the row cursor. This value may be passed to a subsequent call to `mysql_row_seek()`.

Errors

None.

27.7.7.63 mysql_row_tell()

```
MYSQL_ROW_OFFSET mysql_row_tell(MYSQL_RES *result)
```

Description

Returns the current position of the row cursor for the last `mysql_fetch_row()`. This value can be used as an argument to `mysql_row_seek()`.

Use `mysql_row_tell()` only after `mysql_store_result()`, not after `mysql_use_result()`.

Return Values

The current offset of the row cursor.

Errors

None.

27.7.7.64 mysql_select_db()

```
int mysql_select_db(MYSQL *mysql, const char *db)
```

Description

Causes the database specified by `db` to become the default (current) database on the connection specified by `mysql`. In subsequent queries, this database is the default for table references that include no explicit database specifier.

`mysql_select_db()` fails unless the connected user can be authenticated as having permission to use the database.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.65 mysql_server_init()

```
int mysql_server_init(int argc, char **argv, char **groups)
```

Description

This function initializes the MySQL library, which must be done before you call any other MySQL function. However, `mysql_server_init()` is deprecated and you should call `mysql_library_init()` instead. See Section 27.7.7.41, “mysql_library_init()”.

Return Values

Zero for success. Nonzero if an error occurred.

27.7.7.66 mysql_server_end()

```
void mysql_server_end(void)
```

Description

This function finalizes the MySQL library, which should be done when you are done using the library. However, `mysql_server_end()` is deprecated and `mysql_library_end()` should be used instead. See Section 27.7.7.40, “mysql_library_end()”.

Return Values

None.

27.7.7.67 mysql_session_track_get_first()

```
int mysql_session_track_get_first(MYSQL *mysql, enum enum_session_state_type
type, const char **data, size_t *length)
```

Description

This function fetches the first session state-change information received from the server.

To control notification for changes to session state, use the `session_track_state_change`, `session_track_schema`, `session_track_system_variables`, and `session_track_gtids` system variables (see Section 5.1.5, “Server System Variables”).

The function parameters are used as follows. These descriptions also apply to `mysql_session_track_get_next()`, which takes the same parameters.

- `mysql`: The connection handler.
- `type`: The type of information to retrieve. Permitted values for this parameter are the members of the `enum_session_state_type` enumeration defined in `mysql_com.h`:

  ```
  enum enum_session_state_type
  {
    SESSION_TRACK_SYSTEM_VARIABLES, /* Session system variables */
    SESSION_TRACK_SCHEMA,           /* Current schema */
    SESSION_TRACK_STATE_CHANGE      /* track session state changes */
    SESSION_TRACK_GTIDS,            /* track GTIDs/*
  };
  ```

  To make it easy to loop over all possible types of session information, the `SESSION_TRACK_BEGIN` and `SESSION_TRACK_END` macros are defined to be equal to the first and last members of the `enum_session_state_type` enumeration. The example code shown later in this section demonstrates this technique.

- `data`: The address of a `const char *` variable. Following a successful call, this variable points to the returned data, which should be considered read only.
- `length`: The address of a `size_t` variable. Following a successful call, this variable contains the length of the data pointed to by the `data` parameter.

Following a successful call, interpret the `data` and `length` values according to the `type` value, as follows:

- `SESSION_TRACK_SCHEMA`: `data` is the new default schema name and `length` is the name length.
- `SESSION_TRACK_SYSTEM_VARIABLES`: When a session system variable changes, two values per variable are returned (in separate calls). For the first call, `data` is the variable name and `length` is the name length. For the second call, `data` is the variable value and `length` is the value length. value Both data values are represented as strings.
- `SESSION_TRACK_STATE_CHANGE`: `data` is a byte containing a boolean flag that indicates whether session state changes occurred and `length` should be 1. The flag is represented as an ASCII value, not a binary (for example, `'1'`, not `0x01`).

- `SESSION_TRACK_GTIDS`: `data` contains the GTID set information and `length` is its length.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

None.

Example

The following example shows how to call `mysql_session_track_get_first()` and `mysql_session_track_get_next()` to retrieve and display all available session state-change information following successful execution of an SQL statement string (represented by `stmt_str`).

```
printf("Execute: %s\n", stmt_str);

if (mysql_query(mysql, stmt_str) != 0)
{
  fprintf(stderr, "Error %u: %s\n",
          mysql_errno(mysql), mysql_error(mysql));
  return;
}

MYSQL_RES *result = mysql_store_result(mysql);
if (result) /* there is a result set to fetch */
{
  /* ... process rows here ... */
  printf("Number of rows returned: %lu\n",
         (unsigned long) mysql_num_rows(result));
  mysql_free_result(result);
}
else        /* there is no result set */
{
  if (mysql_field_count(mysql) == 0)
  {
    printf("Number of rows affected: %lu\n",
           (unsigned long) mysql_affected_rows(mysql));
  }
  else      /* an error occurred */
  {
    fprintf(stderr, "Error %u: %s\n",
            mysql_errno(mysql), mysql_error(mysql));
  }
}

/* extract any available session state-change information */
enum enum_session_state_type type;
for (type = SESSION_TRACK_BEGIN; type <= SESSION_TRACK_END; type++)
{
  const char *data;
  size_t length;

  if (mysql_session_track_get_first(mysql, type, &data, &length) == 0)
  {
    printf("Type=%d:\n", type);
    printf("mysql_session_track_get_first() returns: %*.*s\n",
           (int) length, (int) length, data);

    /* check for more data */
    while (mysql_session_track_get_next(mysql, type, &data, &length) == 0)
    {
      printf("mysql_session_track_get_next() returns: %*.*s\n",
             (int) length, (int) length, data);
    }
  }
}
```

27.7.7.68 mysql_session_track_get_next()

```
int mysql_session_track_get_next(MYSQL *mysql, enum enum_session_state_type
type, const char **data, size_t *length)
```

Description

This function fetches session state-change information received from the server, following that retrieved by `mysql_session_track_get_first()`.

Following a successful call to `mysql_session_track_get_first()`, call `mysql_session_track_get_next()` repeatedly until it returns nonzero to indicate no more information is available. The calling sequence for `mysql_session_track_get_next()` is similar to that for `mysql_session_track_get_first()`. For more information and an example that demonstrates both functions, see Section 27.7.7.67, “mysql_session_track_get_first()”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

None.

27.7.7.69 mysql_set_character_set()

```
int mysql_set_character_set(MYSQL *mysql, const char *csname)
```

Description

This function is used to set the default character set for the current connection. The string `csname` specifies a valid character set name. The connection collation becomes the default collation of the character set. This function works like the `SET NAMES` statement, but also sets the value of `mysql->charset`, and thus affects the character set used by `mysql_real_escape_string()`

Return Values

Zero for success. Nonzero if an error occurred.

Example

```
MYSQL mysql;

mysql_init(&mysql);
if (!mysql_real_connect(&mysql,"host","user","passwd","database",0,NULL,0))
{
    fprintf(stderr, "Failed to connect to database: Error: %s\n",
          mysql_error(&mysql));
}

if (!mysql_set_character_set(&mysql, "utf8"))
{
    printf("New client character set: %s\n",
           mysql_character_set_name(&mysql));
}
```

27.7.7.70 mysql_set_local_infile_default()

```
void mysql_set_local_infile_default(MYSQL *mysql);
```

Description

Sets the `LOAD DATA LOCAL INFILE` callback functions to the defaults used internally by the C client library. The library calls this function automatically if `mysql_set_local_infile_handler()` has not been called or does not supply valid functions for each of its callbacks.

Return Values

None.

Errors

None.

27.7.7.71 mysql_set_local_infile_handler()

```
void mysql_set_local_infile_handler(MYSQL *mysql, int (*local_infile_init)
(void **, const char *, void *), int (*local_infile_read)(void *, char *,
unsigned int), void (*local_infile_end)(void *), int (*local_infile_error)
(void *, char*, unsigned int), void *userdata);
```

Description

This function installs callbacks to be used during the execution of `LOAD DATA LOCAL INFILE` statements. It enables application programs to exert control over local (client-side) data file reading. The arguments are the connection handler, a set of pointers to callback functions, and a pointer to a data area that the callbacks can use to share information.

To use `mysql_set_local_infile_handler()`, you must write the following callback functions:

```
int
local_infile_init(void **ptr, const char *filename, void *userdata);
```

The initialization function. This is called once to do any setup necessary, open the data file, allocate data structures, and so forth. The first `void**` argument is a pointer to a pointer. You can set the pointer (that is, `*ptr`) to a value that will be passed to each of the other callbacks (as a `void*`). The callbacks can use this pointed-to value to maintain state information. The `userdata` argument is the same value that is passed to `mysql_set_local_infile_handler()`.

Make the initialization function return zero for success, nonzero for an error.

```
int
local_infile_read(void *ptr, char *buf, unsigned int buf_len);
```

The data-reading function. This is called repeatedly to read the data file. `buf` points to the buffer where the read data is stored, and `buf_len` is the maximum number of bytes that the callback can read and store in the buffer. (It can read fewer bytes, but should not read more.)

The return value is the number of bytes read, or zero when no more data could be read (this indicates EOF). Return a value less than zero if an error occurs.

```
void
local_infile_end(void *ptr)
```

The termination function. This is called once after `local_infile_read()` has returned zero (EOF) or an error. Within this function, deallocate any memory allocated by `local_infile_init()` and perform any other cleanup necessary. It is invoked even if the initialization function returns an error.

```
int
local_infile_error(void *ptr,
                   char *error_msg,
                   unsigned int error_msg_len);
```

The error-handling function. This is called to get a textual error message to return to the user in case any of your other functions returns an error. `error_msg` points to the buffer into which the message is written, and `error_msg_len` is the length of the buffer. Write the message as a null-terminated string, at most `error_msg_len`−1 bytes long.

The return value is the error number.

Typically, the other callbacks store the error message in the data structure pointed to by `ptr`, so that `local_infile_error()` can copy the message from there into `error_msg`.

After calling `mysql_set_local_infile_handler()` in your C code and passing pointers to your callback functions, you can then issue a `LOAD DATA LOCAL INFILE` statement (for example, by using `mysql_query()`). The client library automatically invokes your callbacks. The file name specified in `LOAD DATA LOCAL INFILE` will be passed as the second parameter to the `local_infile_init()` callback.

Return Values

None.

Errors

None.

27.7.7.72 mysql_set_server_option()

```
int mysql_set_server_option(MYSQL *mysql, enum enum_mysql_set_option option)
```

Description

Enables or disables an option for the connection. `option` can have one of the following values.

Option	Description
`MYSQL_OPTION_MULTI_STATEMENTS_ON`	Enable multiple-statement support
`MYSQL_OPTION_MULTI_STATEMENTS_OFF`	Disable multiple-statement support

If you enable multiple-statement support, you should retrieve results from calls to `mysql_query()` or `mysql_real_query()` by using a loop that calls `mysql_next_result()` to determine whether there are more results. For an example, see Section 27.7.20, “C API Support for Multiple Statement Execution”.

Enabling multiple-statement support with `MYSQL_OPTION_MULTI_STATEMENTS_ON` does not have quite the same effect as enabling it by passing the `CLIENT_MULTI_STATEMENTS` flag to `mysql_real_connect()`: `CLIENT_MULTI_STATEMENTS` also enables `CLIENT_MULTI_RESULTS`. If you are using the `CALL` SQL statement in your programs, multiple-result support must be enabled; this means that `MYSQL_OPTION_MULTI_STATEMENTS_ON` by itself is insufficient to permit the use of `CALL`.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `ER_UNKNOWN_COM_ERROR`

The server did not support `mysql_set_server_option()` (which is the case that the server is older than 4.1.1) or the server did not support the option one tried to set.

27.7.7.73 mysql_shutdown()

```
int mysql_shutdown(MYSQL *mysql, enum mysql_enum_shutdown_level
shutdown_level)
```

Description

Note

`mysql_shutdown()` is deprecated and will be removed in a future version of MySQL. Instead, use `mysql_query()` to execute a `SHUTDOWN` statement.

Asks the database server to shut down. The connected user must have the `SHUTDOWN` privilege. MySQL servers support only one type of shutdown; `shutdown_level` must be equal to `SHUTDOWN_DEFAULT`. Dynamically linked executables which have been compiled with older versions of the `libmysqlclient` headers and call `mysql_shutdown()` need to be used with the old `libmysqlclient` dynamic library.

An alternative to `mysql_shutdown()` is to use the `SHUTDOWN` SQL statement.

The shutdown process is described in Section 5.1.12, “The Server Shutdown Process”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.74 mysql_sqlstate()

```
const char *mysql_sqlstate(MYSQL *mysql)
```

Description

Returns a null-terminated string containing the SQLSTATE error code for the most recently executed SQL statement. The error code consists of five characters. `'00000'` means “no error.” The values are specified by ANSI SQL and ODBC. For a list of possible values, see Appendix B, *Errors, Error Codes, and Common Problems*.

SQLSTATE values returned by `mysql_sqlstate()` differ from MySQL-specific error numbers returned by `mysql_errno()`. For example, the `mysql` client program displays errors using the following format, where `1146` is the `mysql_errno()` value and `'42S02'` is the corresponding `mysql_sqlstate()` value:

```
shell> SELECT * FROM no_such_table;
ERROR 1146 (42S02): Table 'test.no_such_table' doesn't exist
```

Not all MySQL error numbers are mapped to SQLSTATE error codes. The value `'HY000'` (general error) is used for unmapped error numbers.

If you call `mysql_sqlstate()` after `mysql_real_connect()` fails, `mysql_sqlstate()` might not return a useful value. For example, this happens if a host is blocked by the server and the connection is closed without any SQLSTATE value being sent to the client.

Return Values

A null-terminated character string containing the SQLSTATE error code.

See Also

See Section 27.7.7.14, “mysql_errno()”, Section 27.7.7.15, “mysql_error()”, and Section 27.7.11.27, “mysql_stmt_sqlstate()”.

27.7.7.75 mysql_ssl_set()

```
bool mysql_ssl_set(MYSQL *mysql, const char *key, const char *cert, const
char *ca, const char *capath, const char *cipher)
```

Description

`mysql_ssl_set()` is used for establishing secure connections using SSL. It must be called before `mysql_real_connect()`.

`mysql_ssl_set()` does nothing unless SSL support is enabled in the client library.

Arguments:

- `mysql`: The connection handler returned from `mysql_init()`.
- `key`: The path name to the key file
- `cert`: The path name to the certificate file
- `ca`: The path name to the certificate authority file
- `capath`: The path name to a directory that contains trusted SSL CA certificates in PEM format
- `cipher`: A list of permissible ciphers to use for SSL encryption

The `mysql` argument must be a valid connection handler. Any unused SSL arguments may be given as `NULL`.

Return Values

This function always returns `0`. If SSL setup is incorrect, a subsequent `mysql_real_connect()` call returns an error when you attempt to connect.

Enforcing an Encrypted Connection

`mysql_ssl_set()` specifies SSL information such as certificate and key files for establishing a secure connection if such connections are available, but does not enforce any requirement that the connection obtained be secure. To require an encrypted connection, the standard MySQL client programs use the following technique, which can also be used by third-party applications:

1. If the `--ssl-mode=REQUIRED` command-line option was specified, turn on SSL by calling `mysql_ssl_set()` to supply the appropriate SSL values. In addition, call `mysql_options()`, passing the `MYSQL_OPT_SSL_MODE` option with a value of `SSL_MODE_REQUIRED`. If the `mysql_options()` call fails, exit with an error.

2. Call `mysql_real_connect()` to connect to the server. The call fails if an encrypted connection cannot be obtained.

27.7.7.76 mysql_stat()

```
const char *mysql_stat(MYSQL *mysql)
```

Description

Returns a character string containing information similar to that provided by the `mysqladmin status` command. This includes uptime in seconds and the number of running threads, questions, reloads, and open tables.

Return Values

A character string describing the server status. `NULL` if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.77 mysql_store_result()

```
MYSQL_RES *mysql_store_result(MYSQL *mysql)
```

Description

After invoking `mysql_query()` or `mysql_real_query()`, you must call `mysql_store_result()` or `mysql_use_result()` for every statement that successfully produces a result set (`SELECT`, `SHOW`, `DESCRIBE`, `EXPLAIN`, `CHECK TABLE`, and so forth). You must also call `mysql_free_result()` after you are done with the result set.

You need not call `mysql_store_result()` or `mysql_use_result()` for other statements, but it does not do any harm or cause any notable performance degradation if you call `mysql_store_result()` in all cases. You can detect whether the statement has a result set by checking whether `mysql_store_result()` returns a nonzero value (more about this later).

If you enable multiple-statement support, you should retrieve results from calls to `mysql_query()` or `mysql_real_query()` by using a loop that calls `mysql_next_result()` to determine whether there are more results. For an example, see Section 27.7.20, “C API Support for Multiple Statement Execution”.

If you want to know whether a statement should return a result set, you can use `mysql_field_count()` to check for this. See Section 27.7.7.22, “mysql_field_count()”.

`mysql_store_result()` reads the entire result of a query to the client, allocates a `MYSQL_RES` structure, and places the result into this structure.

`mysql_store_result()` returns a null pointer if the statement did not return a result set (for example, if it was an `INSERT` statement).

`mysql_store_result()` also returns a null pointer if reading of the result set failed. You can check whether an error occurred by checking whether `mysql_error()` returns a nonempty string, `mysql_errno()` returns nonzero, or `mysql_field_count()` returns zero.

An empty result set is returned if there are no rows returned. (An empty result set differs from a null pointer as a return value.)

After you have called `mysql_store_result()` and gotten back a result that is not a null pointer, you can call `mysql_num_rows()` to find out how many rows are in the result set.

You can call `mysql_fetch_row()` to fetch rows from the result set, or `mysql_row_seek()` and `mysql_row_tell()` to obtain or set the current row position within the result set.

See Section 27.7.18.1, “Why mysql_store_result() Sometimes Returns NULL After mysql_query() Returns Success”.

Return Values

A `MYSQL_RES` result structure with the results. `NULL` (0) if an error occurred.

Errors

`mysql_store_result()` resets `mysql_error()` and `mysql_errno()` if it succeeds.

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.78 mysql_thread_id()

```
unsigned long mysql_thread_id(MYSQL *mysql)
```

Description

Returns the thread ID of the current connection. This value can be used as an argument to `mysql_kill()` to kill the thread.

If the connection is lost and you reconnect with `mysql_ping()`, the thread ID changes. This means you should not get the thread ID and store it for later. You should get it when you need it.

Note

This function does not work correctly if thread IDs become larger than 32 bits, which can occur on some systems. To avoid problems with `mysql_thread_id()`, do not use it. To get the connection ID, execute a `SELECT CONNECTION_ID()` query and retrieve the result.

Return Values

The thread ID of the current connection.

Errors

None.

27.7.7.79 mysql_use_result()

```
MYSQL_RES *mysql_use_result(MYSQL *mysql)
```

Description

After invoking `mysql_query()` or `mysql_real_query()`, you must call `mysql_store_result()` or `mysql_use_result()` for every statement that successfully produces a result set (`SELECT`, `SHOW`, `DESCRIBE`, `EXPLAIN`, `CHECK TABLE`, and so forth). You must also call `mysql_free_result()` after you are done with the result set.

`mysql_use_result()` initiates a result set retrieval but does not actually read the result set into the client like `mysql_store_result()` does. Instead, each row must be retrieved individually by making calls to `mysql_fetch_row()`. This reads the result of a query directly from the server without storing it in a temporary table or local buffer, which is somewhat faster and uses much less memory than `mysql_store_result()`. The client allocates memory only for the current row and a communication buffer that may grow up to `max_allowed_packet` bytes.

On the other hand, you should not use `mysql_use_result()` for locking reads if you are doing a lot of processing for each row on the client side, or if the output is sent to a screen on which the user may type a `^S` (stop scroll). This ties up the server and prevent other threads from updating any tables from which the data is being fetched.

When using `mysql_use_result()`, you must execute `mysql_fetch_row()` until a `NULL` value is returned, otherwise, the unfetched rows are returned as part of the result set for your next query. The C API gives the error `Commands out of sync; you can't run this command now` if you forget to do this!

You may not use `mysql_data_seek()`, `mysql_row_seek()`, `mysql_row_tell()`, `mysql_num_rows()`, or `mysql_affected_rows()` with a result returned from `mysql_use_result()`, nor may you issue other queries until `mysql_use_result()` has finished. (However, after you have fetched all the rows, `mysql_num_rows()` accurately returns the number of rows fetched.)

You must call `mysql_free_result()` once you are done with the result set.

Return Values

A `MYSQL_RES` result structure. `NULL` if an error occurred.

Errors

`mysql_use_result()` resets `mysql_error()` and `mysql_errno()` if it succeeds.

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.7.80 mysql_warning_count()

```
unsigned int mysql_warning_count(MYSQL *mysql)
```

Description

Returns the number of errors, warnings, and notes generated during execution of the previous SQL statement.

Return Values

The warning count.

Errors

None.

27.7.8 C API Prepared Statements

The MySQL client/server protocol provides for the use of prepared statements. This capability uses the `MYSQL_STMT` statement handler data structure returned by the `mysql_stmt_init()` initialization function. Prepared execution is an efficient way to execute a statement more than once. The statement is first parsed to prepare it for execution. Then it is executed one or more times at a later time, using the statement handler returned by the initialization function.

Prepared execution is faster than direct execution for statements executed more than once, primarily because the query is parsed only once. In the case of direct execution, the query is parsed every time it is executed. Prepared execution also can provide a reduction of network traffic because for each execution of the prepared statement, it is necessary only to send the data for the parameters.

Prepared statements might not provide a performance increase in some situations. For best results, test your application both with prepared and nonprepared statements and choose whichever yields best performance.

Another advantage of prepared statements is that it uses a binary protocol that makes data transfer between client and server more efficient.

For a list of SQL statements that can be used as prepared statements, see Section 13.5, “Prepared SQL Statement Syntax”.

Metadata changes to tables or views referred to by prepared statements are detected and cause automatic repreparation of the statement when it is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

27.7.9 C API Prepared Statement Data Structures

Prepared statements use several data structures:

- To obtain a statement handler, pass a `MYSQL` connection handler to `mysql_stmt_init()`, which returns a pointer to a `MYSQL_STMT` data structure. This structure is used for further operations with the statement. To specify the statement to prepare, pass the `MYSQL_STMT` pointer and the statement string to `mysql_stmt_prepare()`.

- To provide input parameters for a prepared statement, set up `MYSQL_BIND` structures and pass them to `mysql_stmt_bind_param()`. To receive output column values, set up `MYSQL_BIND` structures and pass them to `mysql_stmt_bind_result()`.
- The `MYSQL_TIME` structure is used to transfer temporal data in both directions.

The following discussion describes the prepared statement data types in detail. For examples that show how to use them, see Section 27.7.11.10, “mysql_stmt_execute()”, and Section 27.7.11.11, “mysql_stmt_fetch()”.

- `MYSQL_STMT`

  This structure is a handler for a prepared statement. A handler is created by calling `mysql_stmt_init()`, which returns a pointer to a `MYSQL_STMT`. The handler is used for all subsequent operations with the statement until you close it with `mysql_stmt_close()`, at which point the handler becomes invalid and should no longer be used.

  The `MYSQL_STMT` structure has no members intended for application use. Applications should not try to copy a `MYSQL_STMT` structure. There is no guarantee that such a copy will be usable.

  Multiple statement handlers can be associated with a single connection. The limit on the number of handlers depends on the available system resources.

- `MYSQL_BIND`

  This structure is used both for statement input (data values sent to the server) and output (result values returned from the server):

  - For input, use `MYSQL_BIND` structures with `mysql_stmt_bind_param()` to bind parameter data values to buffers for use by `mysql_stmt_execute()`.
  - For output, use `MYSQL_BIND` structures with `mysql_stmt_bind_result()` to bind buffers to result set columns, for use in fetching rows with `mysql_stmt_fetch()`.

  To use a `MYSQL_BIND` structure, zero its contents to initialize it, then set its members appropriately. For example, to declare and initialize an array of three `MYSQL_BIND` structures, use this code:

  ```
  MYSQL_BIND bind[3];
  memset(bind, 0, sizeof(bind));
  ```

  The `MYSQL_BIND` structure contains the following members for use by application programs. For several of the members, the manner of use depends on whether the structure is used for input or output.

  - `enum enum_field_types buffer_type`

    The type of the buffer. This member indicates the data type of the C language variable bound to a statement parameter or result set column. For input, `buffer_type` indicates the type of the variable containing the value to be sent to the server. For output, it indicates the type of the variable into which a value received from the server should be stored. For permissible `buffer_type` values, see Section 27.7.9.1, “C API Prepared Statement Type Codes”.

  - `void *buffer`

    A pointer to the buffer to be used for data transfer. This is the address of a C language variable.

    For input, `buffer` is a pointer to the variable in which you store the data value for a statement parameter. When you call `mysql_stmt_execute()`, MySQL use the value stored in the variable in place of the corresponding parameter marker in the statement (specified with `?` in the statement string).

For output, `buffer` is a pointer to the variable in which to return a result set column value. When you call `mysql_stmt_fetch()`, MySQL stores a column value from the current row of the result set in this variable. You can access the value when the call returns.

To minimize the need for MySQL to perform type conversions between C language values on the client side and SQL values on the server side, use C variables that have types similar to those of the corresponding SQL values:

- For numeric data types, `buffer` should point to a variable of the proper numeric C type. For integer variables (which can be `char` for single-byte values or an integer type for larger values), you should also indicate whether the variable has the `unsigned` attribute by setting the `is_unsigned` member, described later.
- For character (nonbinary) and binary string data types, `buffer` should point to a character buffer.
- For date and time data types, `buffer` should point to a `MYSQL_TIME` structure.

For guidelines about mapping between C types and SQL types and notes about type conversions, see Section 27.7.9.1, “C API Prepared Statement Type Codes”, and Section 27.7.9.2, “C API Prepared Statement Type Conversions”.

- `unsigned long buffer_length`

  The actual size of `*buffer` in bytes. This indicates the maximum amount of data that can be stored in the buffer. For character and binary C data, the `buffer_length` value specifies the length of `*buffer` when used with `mysql_stmt_bind_param()` to specify input values, or the maximum number of output data bytes that can be fetched into the buffer when used with `mysql_stmt_bind_result()`.

- `unsigned long *length`

  A pointer to an `unsigned long` variable that indicates the actual number of bytes of data stored in `*buffer`. `length` is used for character or binary C data.

  For input parameter data binding, set `*length` to indicate the actual length of the parameter value stored in `*buffer`. This is used by `mysql_stmt_execute()`.

  For output value binding, MySQL sets `*length` when you call `mysql_stmt_fetch()`. The `mysql_stmt_fetch()` return value determines how to interpret the length:

  - If the return value is 0, `*length` indicates the actual length of the parameter value.
  - If the return value is `MYSQL_DATA_TRUNCATED`, `*length` indicates the nontruncated length of the parameter value. In this case, the minimum of `*length` and `buffer_length` indicates the actual length of the value.

  `length` is ignored for numeric and temporal data types because the `buffer_type` value determines the length of the data value.

  If you must determine the length of a returned value before fetching it, see Section 27.7.11.11, “mysql_stmt_fetch()”, for some strategies.

- `bool *is_null`

  This member points to a `bool` variable that is true if a value is `NULL`, false if it is not `NULL`. For input, set `*is_null` to true to indicate that you are passing a `NULL` value as a statement parameter.

`is_null` is a *pointer* to a boolean scalar, not a boolean scalar, to provide flexibility in how you specify `NULL` values:

- If your data values are always `NULL`, use `MYSQL_TYPE_NULL` as the `buffer_type` value when you bind the column. The other `MYSQL_BIND` members, including `is_null`, do not matter.
- If your data values are always `NOT NULL`, set `is_null = (bool*) 0`, and set the other members appropriately for the variable you are binding.
- In all other cases, set the other members appropriately and set `is_null` to the address of a `bool` variable. Set that variable's value to true or false appropriately between executions to indicate whether the corresponding data value is `NULL` or `NOT NULL`, respectively.

For output, when you fetch a row, MySQL sets the value pointed to by `is_null` to true or false according to whether the result set column value returned from the statement is or is not `NULL`.

- `bool is_unsigned`

  This member applies for C variables with data types that can be `unsigned` (`char`, `short int`, `int`, `long long int`). Set `is_unsigned` to true if the variable pointed to by `buffer` is `unsigned` and false otherwise. For example, if you bind a `signed char` variable to `buffer`, specify a type code of `MYSQL_TYPE_TINY` and set `is_unsigned` to false. If you bind an `unsigned char` instead, the type code is the same but `is_unsigned` should be true. (For `char`, it is not defined whether it is signed or unsigned, so it is best to be explicit about signedness by using `signed char` or `unsigned char`.)

  `is_unsigned` applies only to the C language variable on the client side. It indicates nothing about the signedness of the corresponding SQL value on the server side. For example, if you use an `int` variable to supply a value for a `BIGINT UNSIGNED` column, `is_unsigned` should be false because `int` is a signed type. If you use an `unsigned int` variable to supply a value for a `BIGINT` column, `is_unsigned` should be true because `unsigned int` is an unsigned type. MySQL performs the proper conversion between signed and unsigned values in both directions, although a warning occurs if truncation results.

- `bool *error`

  For output, set this member to point to a `bool` variable to have truncation information for the parameter stored there after a row fetching operation. When truncation reporting is enabled, `mysql_stmt_fetch()` returns `MYSQL_DATA_TRUNCATED` and `*error` is true in the `MYSQL_BIND` structures for parameters in which truncation occurred. Truncation indicates loss of sign or significant digits, or that a string was too long to fit in a column. Truncation reporting is enabled by default, but can be controlled by calling `mysql_options()` with the `MYSQL_REPORT_DATA_TRUNCATION` option.

- `MYSQL_TIME`

  This structure is used to send and receive `DATE`, `TIME`, `DATETIME`, and `TIMESTAMP` data directly to and from the server. Set the `buffer` member to point to a `MYSQL_TIME` structure, and set the `buffer_type` member of a `MYSQL_BIND` structure to one of the temporal types (`MYSQL_TYPE_TIME`, `MYSQL_TYPE_DATE`, `MYSQL_TYPE_DATETIME`, `MYSQL_TYPE_TIMESTAMP`).

  The `MYSQL_TIME` structure contains the members listed in the following table.

Member	Description
`unsigned int year`	The year
`unsigned int month`	The month of the year
`unsigned int day`	The day of the month

Member	Description
`unsigned int hour`	The hour of the day
`unsigned int minute`	The minute of the hour
`unsigned int second`	The second of the minute
`bool neg`	A boolean flag indicating whether the time is negative
`unsigned long second_part`	The fractional part of the second in microseconds

Only those parts of a `MYSQL_TIME` structure that apply to a given type of temporal value are used. The `year`, `month`, and `day` elements are used for `DATE`, `DATETIME`, and `TIMESTAMP` values. The `hour`, `minute`, and `second` elements are used for `TIME`, `DATETIME`, and `TIMESTAMP` values. See Section 27.7.22, “C API Prepared Statement Handling of Date and Time Values”.

27.7.9.1 C API Prepared Statement Type Codes

The `buffer_type` member of `MYSQL_BIND` structures indicates the data type of the C language variable bound to a statement parameter or result set column. For input, `buffer_type` indicates the type of the variable containing the value to be sent to the server. For output, it indicates the type of the variable into which a value received from the server should be stored.

The following table shows the permissible values for the `buffer_type` member of `MYSQL_BIND` structures for input values sent to the server. The table shows the C variable types that you can use, the corresponding type codes, and the SQL data types for which the supplied value can be used without conversion. Choose the `buffer_type` value according to the data type of the C language variable that you are binding. For the integer types, you should also set the `is_unsigned` member to indicate whether the variable is signed or unsigned.

Input Variable C Type	`buffer_type` Value	SQL Type of Destination Value
`signed char`	`MYSQL_TYPE_TINY`	`TINYINT`
`short int`	`MYSQL_TYPE_SHORT`	`SMALLINT`
`int`	`MYSQL_TYPE_LONG`	`INT`
`long long int`	`MYSQL_TYPE_LONGLONG`	`BIGINT`
`float`	`MYSQL_TYPE_FLOAT`	`FLOAT`
`double`	`MYSQL_TYPE_DOUBLE`	`DOUBLE`
`MYSQL_TIME`	`MYSQL_TYPE_TIME`	`TIME`
`MYSQL_TIME`	`MYSQL_TYPE_DATE`	`DATE`
`MYSQL_TIME`	`MYSQL_TYPE_DATETIME`	`DATETIME`
`MYSQL_TIME`	`MYSQL_TYPE_TIMESTAMP`	`TIMESTAMP`
`char[]`	`MYSQL_TYPE_STRING`	`TEXT`, `CHAR`, `VARCHAR`
`char[]`	`MYSQL_TYPE_BLOB`	`BLOB`, `BINARY`, `VARBINARY`
	`MYSQL_TYPE_NULL`	`NULL`

Use `MYSQL_TYPE_NULL` as indicated in the description for the `is_null` member in Section 27.7.9, “C API Prepared Statement Data Structures”.

For input string data, use `MYSQL_TYPE_STRING` or `MYSQL_TYPE_BLOB` depending on whether the value is a character (nonbinary) or binary string:

- `MYSQL_TYPE_STRING` indicates character input string data. The value is assumed to be in the character set indicated by the `character_set_client` system variable. If the server stores the value into a column with a different character set, it converts the value to that character set.
- `MYSQL_TYPE_BLOB` indicates binary input string data. The value is treated as having the `binary` character set. That is, it is treated as a byte string and no conversion occurs.

The following table shows the permissible values for the `buffer_type` member of `MYSQL_BIND` structures for output values received from the server. The table shows the SQL types of received values, the corresponding type codes that such values have in result set metadata, and the recommended C language data types to bind to the `MYSQL_BIND` structure to receive the SQL values without conversion. Choose the `buffer_type` value according to the data type of the C language variable that you are binding. For the integer types, you should also set the `is_unsigned` member to indicate whether the variable is signed or unsigned.

SQL Type of Received Value	`buffer_type` Value	Output Variable C Type
`TINYINT`	`MYSQL_TYPE_TINY`	`signed char`
`SMALLINT`	`MYSQL_TYPE_SHORT`	`short int`
`MEDIUMINT`	`MYSQL_TYPE_INT24`	`int`
`INT`	`MYSQL_TYPE_LONG`	`int`
`BIGINT`	`MYSQL_TYPE_LONGLONG`	`long long int`
`FLOAT`	`MYSQL_TYPE_FLOAT`	`float`
`DOUBLE`	`MYSQL_TYPE_DOUBLE`	`double`
`DECIMAL`	`MYSQL_TYPE_NEWDECIMAL`	`char[]`
`YEAR`	`MYSQL_TYPE_SHORT`	`short int`
`TIME`	`MYSQL_TYPE_TIME`	`MYSQL_TIME`
`DATE`	`MYSQL_TYPE_DATE`	`MYSQL_TIME`
`DATETIME`	`MYSQL_TYPE_DATETIME`	`MYSQL_TIME`
`TIMESTAMP`	`MYSQL_TYPE_TIMESTAMP`	`MYSQL_TIME`
`CHAR`, `BINARY`	`MYSQL_TYPE_STRING`	`char[]`
`VARCHAR`, `VARBINARY`	`MYSQL_TYPE_VAR_STRING`	`char[]`
`TINYBLOB`, `TINYTEXT`	`MYSQL_TYPE_TINY_BLOB`	`char[]`
`BLOB`, `TEXT`	`MYSQL_TYPE_BLOB`	`char[]`
`MEDIUMBLOB`, `MEDIUMTEXT`	`MYSQL_TYPE_MEDIUM_BLOB`	`char[]`
`LONGBLOB`, `LONGTEXT`	`MYSQL_TYPE_LONG_BLOB`	`char[]`
`BIT`	`MYSQL_TYPE_BIT`	`char[]`

27.7.9.2 C API Prepared Statement Type Conversions

Prepared statements transmit data between the client and server using C language variables on the client side that correspond to SQL values on the server side. If there is a mismatch between the C variable type on the client side and the corresponding SQL value type on the server side, MySQL performs implicit type conversions in both directions.

MySQL knows the type code for the SQL value on the server side. The `buffer_type` value in the `MYSQL_BIND` structure indicates the type code of the C variable that holds the value on the client side. The two codes together tell MySQL what conversion must be performed, if any. Here are some examples:

- If you use `MYSQL_TYPE_LONG` with an `int` variable to pass an integer value to the server that is to be stored into a `FLOAT` column, MySQL converts the value to floating-point format before storing it.
- If you fetch an SQL `MEDIUMINT` column value, but specify a `buffer_type` value of `MYSQL_TYPE_LONGLONG` and use a C variable of type `long long int` as the destination buffer, MySQL converts the `MEDIUMINT` value (which requires less than 8 bytes) for storage into the `long long int` (an 8-byte variable).

- If you fetch a numeric column with a value of 255 into a `char[4]` character array and specify a `buffer_type` value of `MYSQL_TYPE_STRING`, the resulting value in the array is a 4-byte string `'255\0'`.
- MySQL returns `DECIMAL` values as the string representation of the original server-side value, which is why the corresponding C type is `char[]`. For example, `12.345` is returned to the client as `'12.345'`. If you specify `MYSQL_TYPE_NEWDECIMAL` and bind a string buffer to the `MYSQL_BIND` structure, `mysql_stmt_fetch()` stores the value in the buffer as a string without conversion. If instead you specify a numeric variable and type code, `mysql_stmt_fetch()` converts the string-format `DECIMAL` value to numeric form.
- For the `MYSQL_TYPE_BIT` type code, `BIT` values are returned into a string buffer, which is why the corresponding C type is `char[]`. The value represents a bit string that requires interpretation on the client side. To return the value as a type that is easier to deal with, you can cause the value to be cast to integer using either of the following types of expressions:

  ```
  SELECT bit_col + 0 FROM t
  SELECT CAST(bit_col AS UNSIGNED) FROM t
  ```

  To retrieve the value, bind an integer variable large enough to hold the value and specify the appropriate corresponding integer type code.

Before binding variables to the `MYSQL_BIND` structures that are to be used for fetching column values, you can check the type codes for each column of the result set. This might be desirable if you want to determine which variable types would be best to use to avoid type conversions. To get the type codes, call `mysql_stmt_result_metadata()` after executing the prepared statement with `mysql_stmt_execute()`. The metadata provides access to the type codes for the result set as described in Section 27.7.11.23, “mysql_stmt_result_metadata()”, and Section 27.7.5, “C API Data Structures”.

To determine whether output string values in a result set returned from the server contain binary or nonbinary data, check whether the `charsetnr` value of the result set metadata is 63 (see Section 27.7.5, “C API Data Structures”). If so, the character set is `binary`, which indicates binary rather than nonbinary data. This enables you to distinguish `BINARY` from `CHAR`, `VARBINARY` from `VARCHAR`, and the `BLOB` types from the `TEXT` types.

If you cause the `max_length` member of the `MYSQL_FIELD` column metadata structures to be set (by calling `mysql_stmt_attr_set()`), be aware that the `max_length` values for the result set indicate the lengths of the longest string representation of the result values, not the lengths of the binary representation. That is, `max_length` does not necessarily correspond to the size of the buffers needed to fetch the values with the binary protocol used for prepared statements. Choose the size of the buffers according to the types of the variables into which you fetch the values. For example, a `TINYINT` column containing the value -128 might have a `max_length` value of 4. But the binary representation of any `TINYINT` value requires only 1 byte for storage, so you can supply a `signed char` variable in which to store the value and set `is_unsigned` to indicate that values are signed.

Metadata changes to tables or views referred to by prepared statements are detected and cause automatic repreparation of the statement when it is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

27.7.10 C API Prepared Statement Function Overview

The following table summarizes the functions available for prepared statement processing. For greater detail, see the descriptions in Section 27.7.11, “C API Prepared Statement Function Descriptions”.

Function	Description
`mysql_stmt_affected_rows()`	Returns the number of rows changed, deleted, or inserted by prepared `UPDATE`, `DELETE`, or `INSERT` statement
`mysql_stmt_attr_get()`	Gets value of an attribute for a prepared statement

Function	Description
mysql_stmt_attr_set()	Sets an attribute for a prepared statement
mysql_stmt_bind_param()	Associates application data buffers with the parameter markers in a prepared SQL statement
mysql_stmt_bind_result()	Associates application data buffers with columns in a result set
mysql_stmt_close()	Frees memory used by a prepared statement
mysql_stmt_data_seek()	Seeks to an arbitrary row number in a statement result set
mysql_stmt_errno()	Returns the error number for the last statement execution
mysql_stmt_error()	Returns the error message for the last statement execution
mysql_stmt_execute()	Executes a prepared statement
mysql_stmt_fetch()	Fetches the next row of data from a result set and returns data for all bound columns
mysql_stmt_fetch_column()	Fetch data for one column of the current row of a result set
mysql_stmt_field_count()	Returns the number of result columns for the most recent statement
mysql_stmt_free_result()	Free the resources allocated to a statement handler
mysql_stmt_init()	Allocates memory for a MYSQL_STMT structure and initializes it
mysql_stmt_insert_id()	Returns the ID generated for an AUTO_INCREMENT column by a prepared statement
mysql_stmt_next_result()	Returns/initiates the next result in a multiple-result execution
mysql_stmt_num_rows()	Returns the row count from a buffered statement result set
mysql_stmt_param_count()	Returns the number of parameters in a prepared statement
mysql_stmt_param_metadata()	(Return parameter metadata in the form of a result set) This function does nothing
mysql_stmt_prepare()	Prepares an SQL statement string for execution
mysql_stmt_reset()	Resets the statement buffers in the server
mysql_stmt_result_metadata()	Returns prepared statement metadata in the form of a result set
mysql_stmt_row_seek()	Seeks to a row offset in a statement result set, using value returned from mysql_stmt_row_tell()
mysql_stmt_row_tell()	Returns the statement row cursor position
mysql_stmt_send_long_data()	Sends long data in chunks to server
mysql_stmt_sqlstate()	Returns the SQLSTATE error code for the last statement execution
mysql_stmt_store_result()	Retrieves a complete result set to the client

Call mysql_stmt_init() to create a statement handler, then mysql_stmt_prepare() to prepare the statement string, mysql_stmt_bind_param() to supply the parameter data, and mysql_stmt_execute() to execute the statement. You can repeat the mysql_stmt_execute() by changing parameter values in the respective buffers supplied through mysql_stmt_bind_param().

You can send text or binary data in chunks to server using mysql_stmt_send_long_data(). See Section 27.7.11.26, “mysql_stmt_send_long_data()”.

If the statement is a SELECT or any other statement that produces a result set, mysql_stmt_prepare() also returns the result set metadata information in the form of a MYSQL_RES result set through mysql_stmt_result_metadata().

You can supply the result buffers using mysql_stmt_bind_result(), so that the mysql_stmt_fetch() automatically returns data to these buffers. This is row-by-row fetching.

When statement execution has been completed, close the statement handler using `mysql_stmt_close()` so that all resources associated with it can be freed. At that point the handler becomes invalid and should no longer be used.

If you obtained a `SELECT` statement's result set metadata by calling `mysql_stmt_result_metadata()`, you should also free the metadata using `mysql_free_result()`.

Execution Steps

To prepare and execute a statement, an application follows these steps:

1. Create a prepared statement handler with `mysql_stmt_init()`. To prepare the statement on the server, call `mysql_stmt_prepare()` and pass it a string containing the SQL statement.
2. If the statement will produce a result set, call `mysql_stmt_result_metadata()` to obtain the result set metadata. This metadata is itself in the form of result set, albeit a separate one from the one that contains the rows returned by the query. The metadata result set indicates how many columns are in the result and contains information about each column.
3. Set the values of any parameters using `mysql_stmt_bind_param()`. All parameters must be set. Otherwise, statement execution returns an error or produces unexpected results.
4. Call `mysql_stmt_execute()` to execute the statement.
5. If the statement produces a result set, bind the data buffers to use for retrieving the row values by calling `mysql_stmt_bind_result()`.
6. Fetch the data into the buffers row by row by calling `mysql_stmt_fetch()` repeatedly until no more rows are found.
7. Repeat steps 3 through 6 as necessary, by changing the parameter values and re-executing the statement.

When `mysql_stmt_prepare()` is called, the MySQL client/server protocol performs these actions:

- The server parses the statement and sends the okay status back to the client by assigning a statement ID. It also sends total number of parameters, a column count, and its metadata if it is a result set oriented statement. All syntax and semantics of the statement are checked by the server during this call.
- The client uses this statement ID for the further operations, so that the server can identify the statement from among its pool of statements.

When `mysql_stmt_execute()` is called, the MySQL client/server protocol performs these actions:

- The client uses the statement handler and sends the parameter data to the server.
- The server identifies the statement using the ID provided by the client, replaces the parameter markers with the newly supplied data, and executes the statement. If the statement produces a result set, the server sends the data back to the client. Otherwise, it sends an okay status and the number of rows changed, deleted, or inserted.

When `mysql_stmt_fetch()` is called, the MySQL client/server protocol performs these actions:

- The client reads the data from the current row of the result set and places it into the application data buffers by doing the necessary conversions. If the application buffer type is same as that of the field type returned from the server, the conversions are straightforward.

If an error occurs, you can get the statement error number, error message, and SQLSTATE code using `mysql_stmt_errno()`, `mysql_stmt_error()`, and `mysql_stmt_sqlstate()`, respectively.

Prepared Statement Logging

For prepared statements that are executed with the `mysql_stmt_prepare()` and `mysql_stmt_execute()` C API functions, the server writes `Prepare` and `Execute` lines to the general query log so that you can tell when statements are prepared and executed.

Suppose that you prepare and execute a statement as follows:

1. Call `mysql_stmt_prepare()` to prepare the statement string `"SELECT ?"`.
2. Call `mysql_stmt_bind_param()` to bind the value `3` to the parameter in the prepared statement.
3. Call `mysql_stmt_execute()` to execute the prepared statement.

As a result of the preceding calls, the server writes the following lines to the general query log:

```
Prepare  [1] SELECT ?
Execute  [1] SELECT 3
```

Each `Prepare` and `Execute` line in the log is tagged with a `[N]` statement identifier so that you can keep track of which prepared statement is being logged. *`N`* is a positive integer. If there are multiple prepared statements active simultaneously for the client, *`N`* may be greater than 1. Each `Execute` lines shows a prepared statement after substitution of data values for `?` parameters.

27.7.11 C API Prepared Statement Function Descriptions

To prepare and execute queries, use the functions described in detail in the following sections.

All functions that operate with a `MYSQL_STMT` structure begin with the prefix `mysql_stmt_`.

To create a `MYSQL_STMT` handler, use the `mysql_stmt_init()` function.

27.7.11.1 mysql_stmt_affected_rows()

```
my_ulonglong mysql_stmt_affected_rows(MYSQL_STMT *stmt)
```

Description

`mysql_stmt_affected_rows()` may be called immediately after executing a statement with `mysql_stmt_execute()`. It is like `mysql_affected_rows()` but for prepared statements. For a description of what the affected-rows value returned by this function means, See Section 27.7.7.1, “mysql_affected_rows()”.

Errors

None.

Example

See the Example in Section 27.7.11.10, “mysql_stmt_execute()”.

27.7.11.2 mysql_stmt_attr_get()

```
bool mysql_stmt_attr_get(MYSQL_STMT *stmt, enum enum_stmt_attr_type option,
void *arg)
```

Description

Can be used to get the current value for a statement attribute.

The `option` argument is the option that you want to get; the `arg` should point to a variable that should contain the option value. If the option is an integer, `arg` should point to the value of the integer.

See Section 27.7.11.3, “mysql_stmt_attr_set()”, for a list of options and option types.

Return Values

Zero for success. Nonzero if `option` is unknown.

Errors

None.

27.7.11.3 mysql_stmt_attr_set()

```
bool mysql_stmt_attr_set(MYSQL_STMT *stmt, enum enum_stmt_attr_type option,
const void *arg)
```

Description

Can be used to affect behavior for a prepared statement. This function may be called multiple times to set several options.

The `option` argument is the option that you want to set. The `arg` argument is the value for the option. `arg` should point to a variable that is set to the desired attribute value. The variable type is as indicated in the following table.

The following table shows the possible `option` values.

Option	Argument Type	Function
`STMT_ATTR_UPDATE_MAX_LENGTH`	`bool *`	If set to 1, causes `mysql_stmt_store_result()` to update the metadata `MYSQL_FIELD->max_length` value.
`STMT_ATTR_CURSOR_TYPE`	`unsigned long *`	Type of cursor to open for statement when `mysql_stmt_execute()` is invoked. `*arg` can be `CURSOR_TYPE_NO_CURSOR` (the default) or `CURSOR_TYPE_READ_ONLY`.
`STMT_ATTR_PREFETCH_ROWS`	`unsigned long *`	Number of rows to fetch from server at a time when using a cursor. `*arg` can be in the range from 1 to the maximum value of `unsigned long`. The default is 1.

If you use the `STMT_ATTR_CURSOR_TYPE` option with `CURSOR_TYPE_READ_ONLY`, a cursor is opened for the statement when you invoke `mysql_stmt_execute()`. If there is already an open cursor from a previous `mysql_stmt_execute()` call, it closes the cursor before opening a new one. `mysql_stmt_reset()` also closes any open cursor before preparing the statement for re-execution. `mysql_stmt_free_result()` closes any open cursor.

If you open a cursor for a prepared statement, `mysql_stmt_store_result()` is unnecessary, because that function causes the result set to be buffered on the client side.

Return Values

Zero for success. Nonzero if `option` is unknown.

Errors

None.

Example

The following example opens a cursor for a prepared statement and sets the number of rows to fetch at a time to 5:

```
MYSQL_STMT *stmt;
int rc;
unsigned long type;
unsigned long prefetch_rows = 5;

stmt = mysql_stmt_init(mysql);
type = (unsigned long) CURSOR_TYPE_READ_ONLY;
rc = mysql_stmt_attr_set(stmt, STMT_ATTR_CURSOR_TYPE, (void*) &type);
/* ... check return value ... */
rc = mysql_stmt_attr_set(stmt, STMT_ATTR_PREFETCH_ROWS,
                         (void*) &prefetch_rows);
/* ... check return value ... */
```

27.7.11.4 mysql_stmt_bind_param()

```
bool mysql_stmt_bind_param(MYSQL_STMT *stmt, MYSQL_BIND *bind)
```

Description

`mysql_stmt_bind_param()` is used to bind input data for the parameter markers in the SQL statement that was passed to `mysql_stmt_prepare()`. It uses `MYSQL_BIND` structures to supply the data. `bind` is the address of an array of `MYSQL_BIND` structures. The client library expects the array to contain one element for each `?` parameter marker that is present in the query.

Suppose that you prepare the following statement:

```
INSERT INTO mytbl VALUES(?,?,?)
```

When you bind the parameters, the array of `MYSQL_BIND` structures must contain three elements, and can be declared like this:

```
MYSQL_BIND bind[3];
```

Section 27.7.9, “C API Prepared Statement Data Structures”, describes the members of each `MYSQL_BIND` element and how they should be set to provide input values.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_UNSUPPORTED_PARAM_TYPE`

  The conversion is not supported. Possibly the `buffer_type` value is invalid or is not one of the supported types.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

See the Example in Section 27.7.11.10, “mysql_stmt_execute()”.

27.7.11.5 mysql_stmt_bind_result()

```
bool mysql_stmt_bind_result(MYSQL_STMT *stmt, MYSQL_BIND *bind)
```

Description

`mysql_stmt_bind_result()` is used to associate (that is, bind) output columns in the result set to data buffers and length buffers. When `mysql_stmt_fetch()` is called to fetch data, the MySQL client/server protocol places the data for the bound columns into the specified buffers.

All columns must be bound to buffers prior to calling `mysql_stmt_fetch()`. `bind` is the address of an array of `MYSQL_BIND` structures. The client library expects the array to contain one element for each column of the result set. If you do not bind columns to `MYSQL_BIND` structures, `mysql_stmt_fetch()` simply ignores the data fetch. The buffers should be large enough to hold the data values, because the protocol does not return data values in chunks.

A column can be bound or rebound at any time, even after a result set has been partially retrieved. The new binding takes effect the next time `mysql_stmt_fetch()` is called. Suppose that an application binds the columns in a result set and calls `mysql_stmt_fetch()`. The client/server protocol returns data in the bound buffers. Then suppose that the application binds the columns to a different set of buffers. The protocol places data into the newly bound buffers when the next call to `mysql_stmt_fetch()` occurs.

To bind a column, an application calls `mysql_stmt_bind_result()` and passes the type, address, and length of the output buffer into which the value should be stored. Section 27.7.9, “C API Prepared Statement Data Structures”, describes the members of each `MYSQL_BIND` element and how they should be set to receive output values.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_UNSUPPORTED_PARAM_TYPE`

  The conversion is not supported. Possibly the `buffer_type` value is invalid or is not one of the supported types.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

See the Example in Section 27.7.11.11, “mysql_stmt_fetch()”.

27.7.11.6 mysql_stmt_close()

```
bool mysql_stmt_close(MYSQL_STMT *stmt)
```

Description

Closes the prepared statement. `mysql_stmt_close()` also deallocates the statement handler pointed to by `stmt`, which at that point becomes invalid and should no longer be used. For a failed `mysql_stmt_close()` call, do not call `mysql_stmt_error()`, or `mysql_stmt_errno()`, or `mysql_stmt_sqlstate()` to obtain error information because `mysql_stmt_close()` makes the statement handler invalid. Call `mysql_error()`, `mysql_errno()`, or `mysql_sqlstate()` instead.

If the current statement has pending or unread results, this function cancels them so that the next query can be executed.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

See the Example in Section 27.7.11.10, “mysql_stmt_execute()”.

27.7.11.7 mysql_stmt_data_seek()

```
void mysql_stmt_data_seek(MYSQL_STMT *stmt, my_ulonglong offset)
```

Description

Seeks to an arbitrary row in a statement result set. The `offset` value is a row number and should be in the range from `0` to `mysql_stmt_num_rows(stmt)-1`.

This function requires that the statement result set structure contains the entire result of the last executed query, so `mysql_stmt_data_seek()` may be used only in conjunction with `mysql_stmt_store_result()`.

Return Values

None.

Errors

None.

27.7.11.8 mysql_stmt_errno()

```
unsigned int mysql_stmt_errno(MYSQL_STMT *stmt)
```

Description

For the statement specified by `stmt`, `mysql_stmt_errno()` returns the error code for the most recently invoked statement API function that can succeed or fail. A return value of zero means that no error occurred. Client error message numbers are listed in the MySQL `errmsg.h` header file. Server error message numbers are listed in `mysqld_error.h`. Errors also are listed at Appendix B, *Errors, Error Codes, and Common Problems*.

If the failed statement API function was `mysql_stmt_close()`, do not call or `mysql_stmt_errno()` to obtain error information because `mysql_stmt_close()` makes the statement handler invalid. Call `mysql_errno()` instead.

Return Values

An error code value. Zero if no error occurred.

Errors

None.

27.7.11.9 mysql_stmt_error()

```
const char *mysql_stmt_error(MYSQL_STMT *stmt)
```

Description

For the statement specified by `stmt`, `mysql_stmt_error()` returns a null-terminated string containing the error message for the most recently invoked statement API function that can succeed or fail. An empty string (`""`) is returned if no error occurred. Either of these two tests can be used to check for an error:

```
if(*mysql_stmt_errno(stmt))
{
  // an error occurred
}

if (mysql_stmt_error(stmt)[0])
{
  // an error occurred
}
```

If the failed statement API function was `mysql_stmt_close()`, do not call `mysql_stmt_error()` to obtain error information because `mysql_stmt_close()` makes the statement handler invalid. Call `mysql_error()` instead.

The language of the client error messages may be changed by recompiling the MySQL client library. You can choose error messages in several different languages.

Return Values

A character string that describes the error. An empty string if no error occurred.

Errors

None.

27.7.11.10 mysql_stmt_execute()

```
int mysql_stmt_execute(MYSQL_STMT *stmt)
```

Description

`mysql_stmt_execute()` executes the prepared query associated with the statement handler. The currently bound parameter marker values are sent to server during this call, and the server replaces the markers with this newly supplied data.

Statement processing following `mysql_stmt_execute()` depends on the type of statement:

- For an `UPDATE`, `DELETE`, or `INSERT`, the number of changed, deleted, or inserted rows can be found by calling `mysql_stmt_affected_rows()`.
- For a statement such as `SELECT` that generates a result set, you must call `mysql_stmt_fetch()` to fetch the data prior to calling any other functions that result in query processing. For more information on how to fetch the results, refer to Section 27.7.11.11, “mysql_stmt_fetch()”.

  Do not following invocation of `mysql_stmt_execute()` with a call to `mysql_store_result()` or `mysql_use_result()`. Those functions are not intended for processing results from prepared statements.

For statements that generate a result set, you can request that `mysql_stmt_execute()` open a cursor for the statement by calling `mysql_stmt_attr_set()` before executing the statement. If you execute a statement multiple times, `mysql_stmt_execute()` closes any open cursor before opening a new one.

Metadata changes to tables or views referred to by prepared statements are detected and cause automatic repreparation of the statement when it is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

The following example demonstrates how to create and populate a table using `mysql_stmt_init()`, `mysql_stmt_prepare()`, `mysql_stmt_param_count()`, `mysql_stmt_bind_param()`, `mysql_stmt_execute()`, and `mysql_stmt_affected_rows()`. The `mysql` variable is assumed to be a valid connection handler. For an example that shows how to retrieve data, see Section 27.7.11.11, “mysql_stmt_fetch()”.

```
#define STRING_SIZE 50

#define DROP_SAMPLE_TABLE "DROP TABLE IF EXISTS test_table"
#define CREATE_SAMPLE_TABLE "CREATE TABLE test_table(col1 INT,\
                                                     col2 VARCHAR(40),\
                                                     col3 SMALLINT,\
                                                     col4 TIMESTAMP)"
#define INSERT_SAMPLE "INSERT INTO \
                       test_table(col1,col2,col3) \
                       VALUES(?,?,?)"

MYSQL_STMT    *stmt;
MYSQL_BIND    bind[3];
my_ulonglong  affected_rows;
int           param_count;
short         small_data;
int           int_data;
char          str_data[STRING_SIZE];
unsigned long str_length;
bool          is_null;

if (mysql_query(mysql, DROP_SAMPLE_TABLE))
{
  fprintf(stderr, " DROP TABLE failed\n");
  fprintf(stderr, " %s\n", mysql_error(mysql));
  exit(0);
}

if (mysql_query(mysql, CREATE_SAMPLE_TABLE))
```

```
{
  fprintf(stderr, " CREATE TABLE failed\n");
  fprintf(stderr, " %s\n", mysql_error(mysql));
  exit(0);
}

/* Prepare an INSERT query with 3 parameters */
/* (the TIMESTAMP column is not named; the server */
/*  sets it to the current date and time) */
stmt = mysql_stmt_init(mysql);
if (!stmt)
{
  fprintf(stderr, " mysql_stmt_init(), out of memory\n");
  exit(0);
}
if (mysql_stmt_prepare(stmt, INSERT_SAMPLE, strlen(INSERT_SAMPLE)))
{
  fprintf(stderr, " mysql_stmt_prepare(), INSERT failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}
fprintf(stdout, " prepare, INSERT successful\n");

/* Get the parameter count from the statement */
param_count= mysql_stmt_param_count(stmt);
fprintf(stdout, " total parameters in INSERT: %d\n", param_count);

if (param_count != 3) /* validate parameter count */
{
  fprintf(stderr, " invalid parameter count returned by MySQL\n");
  exit(0);
}

/* Bind the data for all 3 parameters */

memset(bind, 0, sizeof(bind));

/* INTEGER PARAM */
/* This is a number type, so there is no need
   to specify buffer_length */
bind[0].buffer_type= MYSQL_TYPE_LONG;
bind[0].buffer= (char *)&int_data;
bind[0].is_null= 0;
bind[0].length= 0;

/* STRING PARAM */
bind[1].buffer_type= MYSQL_TYPE_STRING;
bind[1].buffer= (char *)str_data;
bind[1].buffer_length= STRING_SIZE;
bind[1].is_null= 0;
bind[1].length= &str_length;

/* SMALLINT PARAM */
bind[2].buffer_type= MYSQL_TYPE_SHORT;
bind[2].buffer= (char *)&small_data;
bind[2].is_null= &is_null;
bind[2].length= 0;

/* Bind the buffers */
if (mysql_stmt_bind_param(stmt, bind))
{
  fprintf(stderr, " mysql_stmt_bind_param() failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Specify the data values for the first row */
int_data= 10;             /* integer */
strncpy(str_data, "MySQL", STRING_SIZE); /* string  */
str_length= strlen(str_data);

/* INSERT SMALLINT data as NULL */
```

```
is_null= 1;

/* Execute the INSERT statement - 1*/
if (mysql_stmt_execute(stmt))
{
  fprintf(stderr, " mysql_stmt_execute(), 1 failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Get the number of affected rows */
affected_rows= mysql_stmt_affected_rows(stmt);
fprintf(stdout, " total affected rows(insert 1): %lu\n",
                (unsigned long) affected_rows);

if (affected_rows != 1) /* validate affected rows */
{
  fprintf(stderr, " invalid affected rows by MySQL\n");
  exit(0);
}

/* Specify data values for second row,
   then re-execute the statement */
int_data= 1000;
strncpy(str_data, "
        The most popular Open Source database",
        STRING_SIZE);
str_length= strlen(str_data);
small_data= 1000;         /* smallint */
is_null= 0;               /* reset */

/* Execute the INSERT statement - 2*/
if (mysql_stmt_execute(stmt))
{
  fprintf(stderr, " mysql_stmt_execute, 2 failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Get the total rows affected */
affected_rows= mysql_stmt_affected_rows(stmt);
fprintf(stdout, " total affected rows(insert 2): %lu\n",
                (unsigned long) affected_rows);

if (affected_rows != 1) /* validate affected rows */
{
  fprintf(stderr, " invalid affected rows by MySQL\n");
  exit(0);
}

/* Close the statement */
if (mysql_stmt_close(stmt))
{
  /* mysql_stmt_close() invalidates stmt, so call          */
  /* mysql_error(mysql) rather than mysql_stmt_error(stmt) */
  fprintf(stderr, " failed while closing the statement\n");
  fprintf(stderr, " %s\n", mysql_error(mysql));
  exit(0);
}
```

Note

For complete examples on the use of prepared statement functions, refer to the file `tests/mysql_client_test.c`. This file can be obtained from a MySQL source distribution or from the source repository (see Section 2.8, “Installing MySQL from Source”).

27.7.11.11 mysql_stmt_fetch()

```
int mysql_stmt_fetch(MYSQL_STMT *stmt)
```

Description

`mysql_stmt_fetch()` returns the next row in the result set. It can be called only while the result set exists; that is, after a call to `mysql_stmt_execute()` for a statement such as `SELECT` that produces a result set.

`mysql_stmt_fetch()` returns row data using the buffers bound by `mysql_stmt_bind_result()`. It returns the data in those buffers for all the columns in the current row set and the lengths are returned to the `length` pointer. All columns must be bound by the application before it calls `mysql_stmt_fetch()`.

By default, result sets are fetched unbuffered a row at a time from the server. To buffer the entire result set on the client, call `mysql_stmt_store_result()` after binding the data buffers and before calling `mysql_stmt_fetch()`.

If a fetched data value is a `NULL` value, the `*is_null` value of the corresponding `MYSQL_BIND` structure contains TRUE (1). Otherwise, the data and its length are returned in the `*buffer` and `*length` elements based on the buffer type specified by the application. Each numeric and temporal type has a fixed length, as listed in the following table. The length of the string types depends on the length of the actual data value, as indicated by `data_length`.

Type	Length
`MYSQL_TYPE_TINY`	1
`MYSQL_TYPE_SHORT`	2
`MYSQL_TYPE_LONG`	4
`MYSQL_TYPE_LONGLONG`	8
`MYSQL_TYPE_FLOAT`	4
`MYSQL_TYPE_DOUBLE`	8
`MYSQL_TYPE_TIME`	`sizeof(MYSQL_TIME)`
`MYSQL_TYPE_DATE`	`sizeof(MYSQL_TIME)`
`MYSQL_TYPE_DATETIME`	`sizeof(MYSQL_TIME)`
`MYSQL_TYPE_STRING`	`data length`
`MYSQL_TYPE_BLOB`	`data_length`

In some cases you might want to determine the length of a column value before fetching it with `mysql_stmt_fetch()`. For example, the value might be a long string or `BLOB` value for which you want to know how much space must be allocated. To accomplish this, you can use these strategies:

- Before invoking `mysql_stmt_fetch()` to retrieve individual rows, pass `STMT_ATTR_UPDATE_MAX_LENGTH` to `mysql_stmt_attr_set()`, then invoke `mysql_stmt_store_result()` to buffer the entire result on the client side. Setting the `STMT_ATTR_UPDATE_MAX_LENGTH` attribute causes the maximal length of column values to be indicated by the `max_length` member of the result set metadata returned by `mysql_stmt_result_metadata()`.

- Invoke `mysql_stmt_fetch()` with a zero-length buffer for the column in question and a pointer in which the real length can be stored. Then use the real length with `mysql_stmt_fetch_column()`.

```
real_length= 0;

bind[0].buffer= 0;
bind[0].buffer_length= 0;
bind[0].length= &real_length
mysql_stmt_bind_result(stmt, bind);

mysql_stmt_fetch(stmt);
if (real_length > 0)
```

```
{
  data= malloc(real_length);
  bind[0].buffer= data;
  bind[0].buffer_length= real_length;
  mysql_stmt_fetch_column(stmt, bind, 0, 0);
}
```

Return Values

Return Value	Description
0	Successful, the data has been fetched to application data buffers.
1	Error occurred. Error code and message can be obtained by calling `mysql_stmt_errno()` and `mysql_stmt_error()`.
`MYSQL_NO_DATA`	No more rows/data exists
`MYSQL_DATA_TRUNCATED`	Data truncation occurred

`MYSQL_DATA_TRUNCATED` is returned when truncation reporting is enabled. To determine which column values were truncated when this value is returned, check the `error` members of the `MYSQL_BIND` structures used for fetching values. Truncation reporting is enabled by default, but can be controlled by calling `mysql_options()` with the `MYSQL_REPORT_DATA_TRUNCATION` option.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

- `CR_UNSUPPORTED_PARAM_TYPE`

  The buffer type is `MYSQL_TYPE_DATE`, `MYSQL_TYPE_TIME`, `MYSQL_TYPE_DATETIME`, or `MYSQL_TYPE_TIMESTAMP`, but the data type is not `DATE`, `TIME`, `DATETIME`, or `TIMESTAMP`.

- All other unsupported conversion errors are returned from `mysql_stmt_bind_result()`.

Example

The following example demonstrates how to fetch data from a table using `mysql_stmt_result_metadata()`, `mysql_stmt_bind_result()`, and `mysql_stmt_fetch()`. (This example expects to retrieve the two rows inserted by the example shown in Section 27.7.11.10, “mysql_stmt_execute()”.) The `mysql` variable is assumed to be a valid connection handler.

```
#define STRING_SIZE 50

#define SELECT_SAMPLE "SELECT col1, col2, col3, col4 \
                       FROM test_table"

MYSQL_STMT    *stmt;
```

```
MYSQL_BIND    bind[4];
MYSQL_RES     *prepare_meta_result;
MYSQL_TIME    ts;
unsigned long length[4];
int           param_count, column_count, row_count;
short         small_data;
int           int_data;
char          str_data[STRING_SIZE];
bool          is_null[4];
bool          error[4];

/* Prepare a SELECT query to fetch data from test_table */
stmt = mysql_stmt_init(mysql);
if (!stmt)
{
  fprintf(stderr, " mysql_stmt_init(), out of memory\n");
  exit(0);
}
if (mysql_stmt_prepare(stmt, SELECT_SAMPLE, strlen(SELECT_SAMPLE)))
{
  fprintf(stderr, " mysql_stmt_prepare(), SELECT failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}
fprintf(stdout, " prepare, SELECT successful\n");

/* Get the parameter count from the statement */
param_count= mysql_stmt_param_count(stmt);
fprintf(stdout, " total parameters in SELECT: %d\n", param_count);

if (param_count != 0) /* validate parameter count */
{
  fprintf(stderr, " invalid parameter count returned by MySQL\n");
  exit(0);
}

/* Fetch result set meta information */
prepare_meta_result = mysql_stmt_result_metadata(stmt);
if (!prepare_meta_result)
{
  fprintf(stderr,
         " mysql_stmt_result_metadata(), \
           returned no meta information\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Get total columns in the query */
column_count= mysql_num_fields(prepare_meta_result);
fprintf(stdout,
        " total columns in SELECT statement: %d\n",
        column_count);

if (column_count != 4) /* validate column count */
{
  fprintf(stderr, " invalid column count returned by MySQL\n");
  exit(0);
}

/* Execute the SELECT query */
if (mysql_stmt_execute(stmt))
{
  fprintf(stderr, " mysql_stmt_execute(), failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Bind the result buffers for all 4 columns before fetching them */

memset(bind, 0, sizeof(bind));

/* INTEGER COLUMN */
```

```
bind[0].buffer_type= MYSQL_TYPE_LONG;
bind[0].buffer= (char *)&int_data;
bind[0].is_null= &is_null[0];
bind[0].length= &length[0];
bind[0].error= &error[0];

/* STRING COLUMN */
bind[1].buffer_type= MYSQL_TYPE_STRING;
bind[1].buffer= (char *)str_data;
bind[1].buffer_length= STRING_SIZE;
bind[1].is_null= &is_null[1];
bind[1].length= &length[1];
bind[1].error= &error[1];

/* SMALLINT COLUMN */
bind[2].buffer_type= MYSQL_TYPE_SHORT;
bind[2].buffer= (char *)&small_data;
bind[2].is_null= &is_null[2];
bind[2].length= &length[2];
bind[2].error= &error[2];

/* TIMESTAMP COLUMN */
bind[3].buffer_type= MYSQL_TYPE_TIMESTAMP;
bind[3].buffer= (char *)&ts;
bind[3].is_null= &is_null[3];
bind[3].length= &length[3];
bind[3].error= &error[3];

/* Bind the result buffers */
if (mysql_stmt_bind_result(stmt, bind))
{
  fprintf(stderr, " mysql_stmt_bind_result() failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Now buffer all results to client (optional step) */
if (mysql_stmt_store_result(stmt))
{
  fprintf(stderr, " mysql_stmt_store_result() failed\n");
  fprintf(stderr, " %s\n", mysql_stmt_error(stmt));
  exit(0);
}

/* Fetch all rows */
row_count= 0;
fprintf(stdout, "Fetching results ...\n");
while (!mysql_stmt_fetch(stmt))
{
  row_count++;
  fprintf(stdout, "  row %d\n", row_count);

  /* column 1 */
  fprintf(stdout, "   column1 (integer)  : ");
  if (is_null[0])
    fprintf(stdout, " NULL\n");
  else
    fprintf(stdout, " %d(%ld)\n", int_data, length[0]);

  /* column 2 */
  fprintf(stdout, "   column2 (string)   : ");
  if (is_null[1])
    fprintf(stdout, " NULL\n");
  else
    fprintf(stdout, " %s(%ld)\n", str_data, length[1]);

  /* column 3 */
  fprintf(stdout, "   column3 (smallint) : ");
  if (is_null[2])
    fprintf(stdout, " NULL\n");
  else
    fprintf(stdout, " %d(%ld)\n", small_data, length[2]);
```

```
  /* column 4 */
  fprintf(stdout, "   column4 (timestamp): ");
  if (is_null[3])
    fprintf(stdout, " NULL\n");
  else
    fprintf(stdout, " %04d-%02d-%02d %02d:%02d:%02d (%ld)\n",
                    ts.year, ts.month, ts.day,
                    ts.hour, ts.minute, ts.second,
                    length[3]);
  fprintf(stdout, "\n");
}

/* Validate rows fetched */
fprintf(stdout, " total rows fetched: %d\n", row_count);
if (row_count != 2)
{
  fprintf(stderr, " MySQL failed to return all rows\n");
  exit(0);
}

/* Free the prepared result metadata */
mysql_free_result(prepare_meta_result);


/* Close the statement */
if (mysql_stmt_close(stmt))
{
  /* mysql_stmt_close() invalidates stmt, so call          */
  /* mysql_error(mysql) rather than mysql_stmt_error(stmt) */
  fprintf(stderr, " failed while closing the statement\n");
  fprintf(stderr, " %s\n", mysql_error(mysql));
  exit(0);
}
```

27.7.11.12 mysql_stmt_fetch_column()

```
int mysql_stmt_fetch_column(MYSQL_STMT *stmt, MYSQL_BIND *bind, unsigned int
column, unsigned long offset)
```

Description

Fetches one column from the current result set row. `bind` provides the buffer where data should be placed. It should be set up the same way as for `mysql_stmt_bind_result()`. `column` indicates which column to fetch. The first column is numbered 0. `offset` is the offset within the data value at which to begin retrieving data. This can be used for fetching the data value in pieces. The beginning of the value is offset 0.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_INVALID_PARAMETER_NO`

  Invalid column number.

- `CR_NO_DATA`

  The end of the result set has already been reached.

27.7.11.13 mysql_stmt_field_count()

```
unsigned int mysql_stmt_field_count(MYSQL_STMT *stmt)
```

Description

Returns the number of columns for the most recent statement for the statement handler. This value is zero for statements such as `INSERT` or `DELETE` that do not produce result sets.

`mysql_stmt_field_count()` can be called after you have prepared a statement by invoking `mysql_stmt_prepare()`.

Return Values

An unsigned integer representing the number of columns in a result set.

Errors

None.

27.7.11.14 mysql_stmt_free_result()

```
bool mysql_stmt_free_result(MYSQL_STMT *stmt)
```

Description

Releases memory associated with the result set produced by execution of the prepared statement. If there is a cursor open for the statement, `mysql_stmt_free_result()` closes it.

Return Values

Zero for success. Nonzero if an error occurred.

27.7.11.15 mysql_stmt_init()

```
MYSQL_STMT *mysql_stmt_init(MYSQL *mysql)
```

Description

Creates and returns a `MYSQL_STMT` handler. The handler should be freed with `mysql_stmt_close()`, at which point the handler becomes invalid and should no longer be used.

See also Section 27.7.9, “C API Prepared Statement Data Structures”, for more information.

Return Values

A pointer to a `MYSQL_STMT` structure in case of success. `NULL` if out of memory.

Errors

- `CR_OUT_OF_MEMORY`

  Out of memory.

27.7.11.16 mysql_stmt_insert_id()

```
my_ulonglong mysql_stmt_insert_id(MYSQL_STMT *stmt)
```

Description

Returns the value generated for an `AUTO_INCREMENT` column by the prepared `INSERT` or `UPDATE` statement. Use this function after you have executed a prepared `INSERT` statement on a table which contains an `AUTO_INCREMENT` field.

See Section 27.7.7.38, “mysql_insert_id()”, for more information.

Return Values

Value for `AUTO_INCREMENT` column which was automatically generated or explicitly set during execution of prepared statement, or value generated by `LAST_INSERT_ID(expr)` function. Return value is undefined if statement does not set `AUTO_INCREMENT` value.

Errors

None.

27.7.11.17 mysql_stmt_next_result()

```
int mysql_stmt_next_result(MYSQL_STMT *mysql)
```

Description

This function is used when you use prepared `CALL` statements to execute stored procedures, which can return multiple result sets. Use a loop that calls `mysql_stmt_next_result()` to determine whether there are more results. If a procedure has `OUT` or `INOUT` parameters, their values will be returned as a single-row result set following any other result sets. The values will appear in the order in which they are declared in the procedure parameter list.

`mysql_stmt_next_result()` returns a status to indicate whether more results exist. If `mysql_stmt_next_result()` returns an error, there are no more results.

Before each call to `mysql_stmt_next_result()`, you must call `mysql_stmt_free_result()` for the current result if it produced a result set (rather than just a result status).

After calling `mysql_stmt_next_result()` the state of the connection is as if you had called `mysql_stmt_execute()`. This means that you can call `mysql_stmt_bind_result()`, `mysql_stmt_affected_rows()`, and so forth.

It is also possible to test whether there are more results by calling `mysql_more_results()`. However, this function does not change the connection state, so if it returns true, you must still call `mysql_stmt_next_result()` to advance to the next result.

For an example that shows how to use `mysql_stmt_next_result()`, see Section 27.7.23, “C API Support for Prepared CALL Statements”.

Return Values

Return Value	Description
0	Successful and there are more results
-1	Successful and there are no more results
>0	An error occurred

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

An unknown error occurred.

27.7.11.18 mysql_stmt_num_rows()

```
my_ulonglong mysql_stmt_num_rows(MYSQL_STMT *stmt)
```

Description

Returns the number of rows in the result set.

The use of `mysql_stmt_num_rows()` depends on whether you used `mysql_stmt_store_result()` to buffer the entire result set in the statement handler. If you use `mysql_stmt_store_result()`, `mysql_stmt_num_rows()` may be called immediately. Otherwise, the row count is unavailable unless you count the rows as you fetch them.

`mysql_stmt_num_rows()` is intended for use with statements that return a result set, such as `SELECT`. For statements such as `INSERT`, `UPDATE`, or `DELETE`, the number of affected rows can be obtained with `mysql_stmt_affected_rows()`.

Return Values

The number of rows in the result set.

Errors

None.

27.7.11.19 mysql_stmt_param_count()

```
unsigned long mysql_stmt_param_count(MYSQL_STMT *stmt)
```

Description

Returns the number of parameter markers present in the prepared statement.

Return Values

An unsigned long integer representing the number of parameters in a statement.

Errors

None.

Example

See the Example in Section 27.7.11.10, “mysql_stmt_execute()”.

27.7.11.20 mysql_stmt_param_metadata()

```
MYSQL_RES *mysql_stmt_param_metadata(MYSQL_STMT *stmt)
```

This function currently does nothing.

27.7.11.21 mysql_stmt_prepare()

```
int mysql_stmt_prepare(MYSQL_STMT *stmt, const char *stmt_str, unsigned long
length)
```

Description

Given the statement handler returned by `mysql_stmt_init()`, prepares the SQL statement pointed to by the string `stmt_str` and returns a status value. The string length should be given by the `length`

argument. The string must consist of a single SQL statement. You should not add a terminating semicolon (`;`) or `\g` to the statement.

The application can include one or more parameter markers in the SQL statement by embedding question mark (`?`) characters into the SQL string at the appropriate positions.

The markers are legal only in certain places in SQL statements. For example, they are permitted in the `VALUES()` list of an `INSERT` statement (to specify column values for a row), or in a comparison with a column in a `WHERE` clause to specify a comparison value. However, they are not permitted for identifiers (such as table or column names), or to specify both operands of a binary operator such as the `=` equal sign. The latter restriction is necessary because it would be impossible to determine the parameter type. In general, parameters are legal only in Data Manipulation Language (DML) statements, and not in Data Definition Language (DDL) statements.

The parameter markers must be bound to application variables using `mysql_stmt_bind_param()` before executing the statement.

Metadata changes to tables or views referred to by prepared statements are detected and cause automatic repreparation of the statement when it is next executed. For more information, see Section 8.10.4, “Caching of Prepared Statements and Stored Programs”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

If the prepare operation was unsuccessful (that is, `mysql_stmt_prepare()` returns nonzero), the error message can be obtained by calling `mysql_stmt_error()`.

Example

See the Example in Section 27.7.11.10, “mysql_stmt_execute()”.

27.7.11.22 mysql_stmt_reset()

```
bool mysql_stmt_reset(MYSQL_STMT *stmt)
```

Description

Resets a prepared statement on client and server to state after prepare. It resets the statement on the server, data sent using `mysql_stmt_send_long_data()`, unbuffered result sets and current errors.

It does not clear bindings or stored result sets. Stored result sets will be cleared when executing the prepared statement (or closing it).

To re-prepare the statement with another query, use `mysql_stmt_prepare()`.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.11.23 mysql_stmt_result_metadata()

```
MYSQL_RES *mysql_stmt_result_metadata(MYSQL_STMT *stmt)
```

Description

If a statement passed to `mysql_stmt_prepare()` is one that produces a result set, `mysql_stmt_result_metadata()` returns the result set metadata in the form of a pointer to a `MYSQL_RES` structure that can be used to process the meta information such as number of fields and individual field information. This result set pointer can be passed as an argument to any of the field-based API functions that process result set metadata, such as:

- `mysql_num_fields()`
- `mysql_fetch_field()`
- `mysql_fetch_field_direct()`
- `mysql_fetch_fields()`
- `mysql_field_count()`
- `mysql_field_seek()`
- `mysql_field_tell()`
- `mysql_free_result()`

The result set structure should be freed when you are done with it, which you can do by passing it to `mysql_free_result()`. This is similar to the way you free a result set obtained from a call to `mysql_store_result()`.

The result set returned by `mysql_stmt_result_metadata()` contains only metadata. It does not contain any row results. The rows are obtained by using the statement handler with `mysql_stmt_fetch()`.

Return Values

A `MYSQL_RES` result structure. `NULL` if no meta information exists for the prepared query.

Errors

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

See the Example in Section 27.7.11.11, “mysql_stmt_fetch()”.

27.7.11.24 mysql_stmt_row_seek()

```
MYSQL_ROW_OFFSET mysql_stmt_row_seek(MYSQL_STMT *stmt, MYSQL_ROW_OFFSET
offset)
```

Description

Sets the row cursor to an arbitrary row in a statement result set. The `offset` value is a row offset that should be a value returned from `mysql_stmt_row_tell()` or from `mysql_stmt_row_seek()`. This value is not a row number; if you want to seek to a row within a result set by number, use `mysql_stmt_data_seek()` instead.

This function requires that the result set structure contains the entire result of the query, so `mysql_stmt_row_seek()` may be used only in conjunction with `mysql_stmt_store_result()`.

Return Values

The previous value of the row cursor. This value may be passed to a subsequent call to `mysql_stmt_row_seek()`.

Errors

None.

27.7.11.25 mysql_stmt_row_tell()

```
MYSQL_ROW_OFFSET mysql_stmt_row_tell(MYSQL_STMT *stmt)
```

Description

Returns the current position of the row cursor for the last `mysql_stmt_fetch()`. This value can be used as an argument to `mysql_stmt_row_seek()`.

You should use `mysql_stmt_row_tell()` only after `mysql_stmt_store_result()`.

Return Values

The current offset of the row cursor.

Errors

None.

27.7.11.26 mysql_stmt_send_long_data()

```
bool mysql_stmt_send_long_data(MYSQL_STMT *stmt, unsigned int
parameter_number, const char *data, unsigned long length)
```

Description

Enables an application to send parameter data to the server in pieces (or “chunks”). Call this function after `mysql_stmt_bind_param()` and before `mysql_stmt_execute()`. It can be called multiple times to send the parts of a character or binary data value for a column, which must be one of the `TEXT` or `BLOB` data types.

`parameter_number` indicates which parameter to associate the data with. Parameters are numbered beginning with 0. `data` is a pointer to a buffer containing data to be sent, and `length` indicates the number of bytes in the buffer.

Note

The next `mysql_stmt_execute()` call ignores the bind buffer for all parameters that have been used with `mysql_stmt_send_long_data()` since last `mysql_stmt_execute()` or `mysql_stmt_reset()`.

If you want to reset/forget the sent data, you can do it with `mysql_stmt_reset()`. See Section 27.7.11.22, “mysql_stmt_reset()”.

The `max_allowed_packet` system variable controls the maximum size of parameter values that can be sent with `mysql_stmt_send_long_data()`.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_INVALID_BUFFER_USE`

  The parameter does not have a string or binary type.

- `CR_INVALID_PARAMETER_NO`

  Invalid parameter number.

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

Example

The following example demonstrates how to send the data for a `TEXT` column in chunks. It inserts the data value `'MySQL - The most popular Open Source database'` into the `text_column` column. The `mysql` variable is assumed to be a valid connection handler.

```
#define INSERT_QUERY "INSERT INTO \
                      test_long_data(text_column) VALUES(?)"

MYSQL_BIND bind[1];
long       length;

stmt = mysql_stmt_init(mysql);
if (!stmt)
{
  fprintf(stderr, " mysql_stmt_init(), out of memory\n");
  exit(0);
}
if (mysql_stmt_prepare(stmt, INSERT_QUERY, strlen(INSERT_QUERY)))
{
  fprintf(stderr, "\n mysql_stmt_prepare(), INSERT failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}
 memset(bind, 0, sizeof(bind));
 bind[0].buffer_type= MYSQL_TYPE_STRING;
 bind[0].length= &length;
 bind[0].is_null= 0;

/* Bind the buffers */
if (mysql_stmt_bind_param(stmt, bind))
{
  fprintf(stderr, "\n param bind failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}

 /* Supply data in chunks to server */
 if (mysql_stmt_send_long_data(stmt,0,"MySQL",5))
{
  fprintf(stderr, "\n send_long_data failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}

 /* Supply the next piece of data */
 if (mysql_stmt_send_long_data(stmt,0,
           " - The most popular Open Source database",40))
{
  fprintf(stderr, "\n send_long_data failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}

 /* Now, execute the query */
 if (mysql_stmt_execute(stmt))
{
  fprintf(stderr, "\n mysql_stmt_execute failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}
```

27.7.11.27 mysql_stmt_sqlstate()

```
const char *mysql_stmt_sqlstate(MYSQL_STMT *stmt)
```

Description

For the statement specified by `stmt`, `mysql_stmt_sqlstate()` returns a null-terminated string containing the SQLSTATE error code for the most recently invoked prepared statement API function that can succeed or fail. The error code consists of five characters. `"00000"` means “no error.” The values are specified by ANSI SQL and ODBC. For a list of possible values, see Appendix B, *Errors, Error Codes, and Common Problems*.

Not all MySQL errors are mapped to SQLSTATE codes. The value `"HY000"` (general error) is used for unmapped errors.

If the failed statement API function was `mysql_stmt_close()`, do not call `mysql_stmt_sqlstate()` to obtain error information because `mysql_stmt_close()` makes the statement handler invalid. Call `mysql_sqlstate()` instead.

Return Values

A null-terminated character string containing the SQLSTATE error code.

27.7.11.28 mysql_stmt_store_result()

```
int mysql_stmt_store_result(MYSQL_STMT *stmt)
```

Description

Result sets are produced by calling `mysql_stmt_execute()` to executed prepared statements for SQL statements such as `SELECT`, `SHOW`, `DESCRIBE`, and `EXPLAIN`. By default, result sets for successfully executed prepared statements are not buffered on the client and `mysql_stmt_fetch()` fetches them one at a time from the server. To cause the complete result set to be buffered on the client, call `mysql_stmt_store_result()` after binding data buffers with `mysql_stmt_bind_result()` and before calling `mysql_stmt_fetch()` to fetch rows. (For an example, see Section 27.7.11.11, “mysql_stmt_fetch()”.)

`mysql_stmt_store_result()` is optional for result set processing, unless you will call `mysql_stmt_data_seek()`, `mysql_stmt_row_seek()`, or `mysql_stmt_row_tell()`. Those functions require a seekable result set.

It is unnecessary to call `mysql_stmt_store_result()` after executing an SQL statement that does not produce a result set, but if you do, it does not harm or cause any notable performance problem. You can detect whether the statement produced a result set by checking if `mysql_stmt_result_metadata()` returns `NULL`. For more information, refer to Section 27.7.11.23, “mysql_stmt_result_metadata()”.

Note

MySQL does not by default calculate `MYSQL_FIELD->max_length` for all columns in `mysql_stmt_store_result()` because calculating this would slow down `mysql_stmt_store_result()` considerably and most applications do not need `max_length`. If you want `max_length` to be updated, you can call `mysql_stmt_attr_set(MYSQL_STMT, STMT_ATTR_UPDATE_MAX_LENGTH, &flag)` to enable this. See Section 27.7.11.3, “mysql_stmt_attr_set()”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_COMMANDS_OUT_OF_SYNC`

  Commands were executed in an improper order.

- `CR_OUT_OF_MEMORY`

  Out of memory.

- `CR_SERVER_GONE_ERROR`

  The MySQL server has gone away.

- `CR_SERVER_LOST`

  The connection to the server was lost during the query.

- `CR_UNKNOWN_ERROR`

  An unknown error occurred.

27.7.12 C API Threaded Function Descriptions

To create a threaded client, use the functions described in the following sections. See also Section 27.7.4.3, “Writing C API Threaded Client Programs”.

27.7.12.1 my_init()

```
void my_init(void)
```

Description

Note

As of MySQL 8.0.2, the `my_init()` function is no longer included in the list of symbols exported from `libmysqlclient`. It need not be called explicitly by client programs because it is called implicitly by other C API initialization functions.

`my_init()` initializes some global variables that MySQL needs. It also calls `mysql_thread_init()` for this thread.

It is necessary for `my_init()` to be called early in the initialization phase of a program's use of the MySQL library. However, `my_init()` is automatically called by `mysql_init()`, `mysql_library_init()`, `mysql_server_init()`, and `mysql_connect()`. If you ensure that your program invokes one of those functions before any other MySQL calls, there is no need to invoke `my_init()` explicitly.

To access the prototype for `my_init()`, your program should include these header files:

```
#include <my_global.h>
#include <my_sys.h>
```

Return Values

None.

27.7.12.2 mysql_thread_end()

```
void mysql_thread_end(void)
```

Description

Call this function before calling `pthread_exit()` to free memory allocated by `mysql_thread_init()`.

`mysql_thread_end()` *is not invoked automatically by the client library*. Before MySQL 5.7.9, it must be called for each `mysql_thread_init()` call to avoid a memory leak. As of MySQL 5.7.9, C API internals were reimplemented to reduce the amount of information allocated by `mysql_thread_init()` that must be freed by `mysql_thread_end()`:

- For release/production builds without debugging support enabled, `mysql_thread_end()` need not be called.
- For debug builds, `mysql_thread_init()` allocates debugging information for the DBUG package (see Section 28.5.3, “The DBUG Package”). `mysql_thread_end()` must be called for each `mysql_thread_init()` call to avoid a memory leak.

Return Values

None.

27.7.12.3 mysql_thread_init()

```
bool mysql_thread_init(void)
```

Description

This function must be called early within each created thread to initialize thread-specific variables. However, you may not necessarily need to invoke it explicitly: calling `mysql_thread_init()` is automatically handled by `mysql_init()`, `mysql_library_init()`, `mysql_server_init()`, and `mysql_connect()`. If you invoke any of those functions, `mysql_thread_init()` is called for you.

Return Values

Zero for success. Nonzero if an error occurred.

27.7.12.4 mysql_thread_safe()

```
unsigned int mysql_thread_safe(void)
```

Description

This function indicates whether the client library is compiled as thread-safe.

Return Values

1 if the client library is thread-safe, 0 otherwise.

27.7.13 C API Client Plugin Functions

This section describes functions used for the client-side plugin API. They enable management of client plugins. For a description of the `st_mysql_client_plugin` structure used by these functions, see Client Plugin Descriptors.

It is unlikely that a client program needs to call the functions in this section. For example, a client that supports the use of authentication plugins normally causes a plugin to be loaded by calling `mysql_options()` to set the `MYSQL_DEFAULT_AUTH` and `MYSQL_PLUGIN_DIR` options:

```
char *plugin_dir = "path_to_plugin_dir";
char *default_auth = "plugin_name";

/* ... process command-line options ... */

mysql_options(&mysql, MYSQL_PLUGIN_DIR, plugin_dir);
mysql_options(&mysql, MYSQL_DEFAULT_AUTH, default_auth);
```

Typically, the program will also accept `--plugin-dir` and `--default-auth` options that enable users to override the default values.

27.7.13.1 mysql_client_find_plugin()

```
struct st_mysql_client_plugin *mysql_client_find_plugin(MYSQL *mysql, const
char *name, int type)
```

Description

Returns a pointer to a loaded plugin, loading the plugin first if necessary. An error occurs if the type is invalid or the plugin cannot be found or loaded.

Specify the parameters as follows:

- `mysql`: A pointer to a `MYSQL` structure. The plugin API does not require a connection to a MySQL server, but this structure must be properly initialized. The structure is used to obtain connection-related information.
- `name`: The plugin name.
- `type`: The plugin type.

Return Values

A pointer to the plugin for success. `NULL` if an error occurred.

Errors

To check for errors, call the `mysql_error()` or `mysql_errno()` function. See Section 27.7.7.15, “mysql_error()”, and Section 27.7.7.14, “mysql_errno()”.

Example

```
MYSQL mysql;
struct st_mysql_client_plugin *p;

if ((p = mysql_client_find_plugin(&mysql, "myplugin",
                                  MYSQL_CLIENT_AUTHENTICATION_PLUGIN, 0)))
{
  printf("Plugin version: %d.%d.%d\n", p->version[0], p->version[1], p->version[2]);
}
```

27.7.13.2 mysql_client_register_plugin()

```
struct st_mysql_client_plugin *mysql_client_register_plugin(MYSQL *mysql,
struct st_mysql_client_plugin *plugin)
```

Description

Adds a plugin structure to the list of loaded plugins. An error occurs if the plugin is already loaded.

Specify the parameters as follows:

- `mysql`: A pointer to a `MYSQL` structure. The plugin API does not require a connection to a MySQL server, but this structure must be properly initialized. The structure is used to obtain connection-related information.
- `plugin`: A pointer to the plugin structure.

Return Values

A pointer to the plugin for success. `NULL` if an error occurred.

Errors

To check for errors, call the `mysql_error()` or `mysql_errno()` function. See Section 27.7.7.15, “mysql_error()”, and Section 27.7.7.14, “mysql_errno()”.

27.7.13.3 mysql_load_plugin()

```
struct st_mysql_client_plugin *mysql_load_plugin(MYSQL *mysql, const char
*name, int type, int argc, ...)
```

Description

Loads a MySQL client plugin, specified by name and type. An error occurs if the type is invalid or the plugin cannot be loaded.

It is not possible to load multiple plugins of the same type. An error occurs if you try to load a plugin of a type already loaded.

Specify the parameters as follows:

- `mysql`: A pointer to a `MYSQL` structure. The plugin API does not require a connection to a MySQL server, but this structure must be properly initialized. The structure is used to obtain connection-related information.
- `name`: The name of the plugin to load.
- `type`: The type of plugin to load, or −1 to disable type checking. If type is not −1, only plugins matching the type are considered for loading.
- `argc`: The number of following arguments (0 if there are none). Interpretation of any following arguments depends on the plugin type.

Another way to cause plugins to be loaded is to set the `LIBMYSQL_PLUGINS` environment variable to a semicolon-separated list of plugin names. For example:

```
shell> export LIBMYSQL_PLUGINS="myplugin1;myplugin2"
```

Plugins named by `LIBMYSQL_PLUGINS` are loaded when the client program calls `mysql_library_init()`. No error is reported if problems occur loading these plugins.

The `LIBMYSQL_PLUGIN_DIR` environment variable can be set to the path name of the directory in which to look for client plugins. This variable is used in two ways:

- During client plugin preloading, the value of the `--plugin-dir` option is not available, so client plugin loading fails unless the plugins are located in the hardwired default directory. If the plugins are located elsewhere, `LIBMYSQL_PLUGIN_DIR` environment variable can be set to the proper directory to enable plugin preloading to succeed.
- For explicit client plugin loading, the `mysql_load_plugin()` and `mysql_load_plugin_v()` C API functions use the `LIBMYSQL_PLUGIN_DIR` value if it exists and the `--plugin-dir` option was not given. If `--plugin-dir` is given, `mysql_load_plugin()` and `mysql_load_plugin_v()` ignore `LIBMYSQL_PLUGIN_DIR`.

Return Values

A pointer to the plugin if it was loaded successfully. `NULL` if an error occurred.

Errors

To check for errors, call the `mysql_error()` or `mysql_errno()` function. See Section 27.7.7.15, “mysql_error()”, and Section 27.7.7.14, “mysql_errno()”.

Example

```
MYSQL mysql;

if(!mysql_load_plugin(&mysql, "myplugin",
                      MYSQL_CLIENT_AUTHENTICATION_PLUGIN, 0))
{
    fprintf(stderr, "Error: %s\n", mysql_error(&mysql));
    exit(-1);
}
```

See Also

See also Section 27.7.13.3, “mysql_load_plugin()”, Section 27.7.7.15, “mysql_error()”, Section 27.7.7.14, “mysql_errno()”.

27.7.13.4 mysql_load_plugin_v()

```
struct st_mysql_client_plugin *mysql_load_plugin_v(MYSQL *mysql, const char
*name, int type, int argc, va_list args)
```

Description

This function is equivalent to `mysql_load_plugin()`, but it accepts a `va_list` instead of a variable list of parameters.

See Also

See also Section 27.7.13.3, “mysql_load_plugin()”.

27.7.13.5 mysql_plugin_options()

```
int mysql_plugin_options(struct st_mysql_client_plugin *plugin, const char
*option, const void *value)
```

Description

Passes an option type and value to a plugin. This function can be called multiple times to set several options. If the plugin does not have an option handler, an error occurs.

Specify the parameters as follows:

- `plugin`: A pointer to the plugin structure.
- `option`: The option to be set.
- `value`: A pointer to the option value.

Return Values

Zero for success, 1 if an error occurred. If the plugin has an option handler, that handler should also return zero for success and 1 if an error occurred.

27.7.14 C API Binary Log Interface

The MySQL client/server protocol includes a client interface for reading a stream of replication events from a MySQL server binary log. This capability uses the `MYSQL_RPL` data structure and a small set of functions to manage communication between a client program and the server from which the binary log is to be read. The following sections describe aspects of this interface in more detail.

27.7.15 C API Binary Log Data Structures

C API functions for processing a replication event stream from a server require a connection handler (a `MYSQL *` pointer) and a pointer to a `MYSQL_RPL` structure that describes the steam of replication events to read from the server binary log. For example:

```
MYSQL *mysql = mysql_real_connect(...);

MYSQL_RPL rpl;

# ... initialize MYSQL_RPL members ...

int result = mysql_binlog_open(mysql, &rpl);
```

This section describes the `MYSQL_RPL` structure members. Connection handlers are described in Section 27.7.5, “C API Data Structures”.

The applicable `MYSQL_RPL` members depend on the binary log operation to be performed:

- Before calling `mysql_binlog_open()`, the caller must set the `MYSQL_RPL` members from `file_name_length` through `flags`. In addition, if `flags` has the `MYSQL_RPL_GTID` flag set, the caller must set the members from `gtid_set_encoded_size` through `gtid_set_arg`.
- After a successful `mysql_binlog_fetch()` call, the caller examines the `size` and `buffer` members.

`MYSQL_RPL` structure member descriptions:

- `file_name_length`

  The length of the name of the binary log file to read. This member is used in conjunction with `file_name`; see the `file_name` description.

- `file_name`

  The name of the binary log file to read:

  - If `file_name` is `NULL`, the client library sets it to the empty string and sets `file_name_length` to 0.
  - If `file_name` is not `NULL`, `file_name_length` must either be the length of the name or 0. If `file_name_length` is 0, the client library sets it to the length of the name, in which case, `file_name` must be given as a null-terminated string.

  To read from the beginning of the binary log without having to know the name of the oldest binary log file, set `file_name` to `NULL` or the empty string, and `start_position` to 4.

- `start_position`

  The position at which to start reading the binary log. The position of the first event in any given binary log file is 4.

- `server_id`

  The server ID to use for identifying to the server from which the binary log is read.

- `flags`

  The union of flags that affect binary log reading, or 0 if no flags are set. These flag values are permitted:

  - `MYSQL_RPL_SKIP_HEARTBEAT`

    Set this flag to cause `mysql_binlog_fetch()` to skip heartbeat events.

  - `MYSQL_RPL_GTID`

    Set this flag to read GTID (global transaction ID) data. If set, you must initialize the `MYSQL_RPL` structure GTID-related members from `gtid_set_encoded_size` to `gtid_set_arg` before calling `mysql_binlog_open()`.

    It is beyond the scope of this documentation to describe in detail how client programs use those GTID-related members. For more information, examine the `mysqlbinlog.cc` source file. For information about GTID-based replication, see Section 18.1.3, “Replication with Global Transaction Identifiers”.

- `gtid_set_encoded_size`

  The size of GTID set data, or 0.

- `fix_gtid_set`

The address of a callback function for `mysql_binlog_open()` to call to fill the command packet GTID set, or `NULL` if there is no such function. The callback function, if used, should have this calling signature:

```
void my_callback(MYSQL_RPL *rpl, unsigned char *packet_gtid_set);
```

- `gtid_set_arg`

  Either a pointer to GTID set data (if `fix_gtid_set` is `NULL`), or a pointer to a value to be made available for use within the callback function (if `fix_gtid_set` is not `NULL`). `gtid_set_arg` is a generic pointer, so it can point to any kind of value (for example, a string, a structure, or a function). Its interpretation within the callback depends on how the callback intends to use it.

- `size`

  After a successful `mysql_binlog_fetch()` call, the size of the returned binary log event. The value is 0 for an EOF event, greater than 0 for a non-EOF event.

- `buffer`

  After a successful `mysql_binlog_fetch()` call, a pointer to the binary log event contents.

27.7.16 C API Binary Log Function Overview

The following table summarizes the functions available for reading a replication event stream from a binary log. For greater detail, see the descriptions in Section 27.7.17, “C API Binary Log Function Descriptions”.

Function	Description
`mysql_binlog_close()`	Close replication event stream
`mysql_binlog_fetch()`	Read event from replication event stream
`mysql_binlog_open()`	Open replication event stream

27.7.17 C API Binary Log Function Descriptions

The following sections provide detailed descriptions of the functions that enable reading the stream of replication events from a MySQL server binary log.

The following simple example program demonstrates the binary log C API functions. Program notes:

- `mysql` is assumed to be a valid connection handler.
- The initial `SET` statement sets the `@master_binlog_checksum` user-defined variable that the server takes as an indication that the client is checksum-aware. This client does nothing with checksums, but without this statement, a server that includes checksums in binary log events will return an error for the first attempt to read an event containing a checksum. The value assigned to the variable is immaterial; what matters is that the variable exist. existence.

```
if (mysql_query(mysql, "SET @master_binlog_checksum='ALL'"))
{
  fprintf(stderr, "mysql_query() failed\n");
  fprintf(stderr, "Error %u: %s\n",
          mysql_errno(mysql), mysql_error(mysql));
  exit(1);
}

MYSQL_RPL rpl;

rpl.file_name_length = 0;
rpl.file_name = NULL;
```

```
rpl.start_position = 4;
rpl.server_id = 0;
rpl.flags = 0;

if (mysql_binlog_open(mysql, &rpl))
{
  fprintf(stderr, "mysql_binlog_open() failed\n");
  fprintf(stderr, "Error %u: %s\n",
          mysql_errno(mysql), mysql_error(mysql));
  exit(1);
}
for (;;)  /* read events until error or EOF */
{
  if (mysql_binlog_fetch(mysql, &rpl))
  {
    fprintf(stderr, "mysql_binlog_fetch() failed\n");
    fprintf(stderr, "Error %u: %s\n",
            mysql_errno(mysql), mysql_error(mysql));
    break;
  }
  if (rpl.size == 0)  /* EOF */
  {
    fprintf(stderr, "EOF event received\n");
    break;
  }
  fprintf(stderr, "Event received of size %lu.\n", rpl.size);
}
mysql_binlog_close(mysql, &rpl);
```

For additional examples that show how to use these functions, look in a MySQL source distribution for these source files:

- `mysqlbinlog.cc` in the `client` directory
- `mysql_client_test.c` in the `testclients` directory

27.7.17.1 mysql_binlog_close()

```
void mysql_binlog_close(MYSQL *mysql, MYSQL_RPL *rpl)
```

Description

Close a replication event stream.

Arguments:

- `mysql`: The connection handler returned from `mysql_init()`. The handler remains open after the `mysql_binlog_close()` call.
- `rpl`: The replication stream structure. After calling `mysql_binlog_close()`, this structure should not be used further without reinitializing it and calling `mysql_binlog_open()` again.

Errors

None.

Example

See Section 27.7.17, “C API Binary Log Function Descriptions”.

27.7.17.2 mysql_binlog_fetch()

```
int mysql_binlog_fetch(MYSQL *mysql, MYSQL_RPL *rpl)
```

Description

Fetch one event from the replication event stream.

Arguments:

- `mysql`: The connection handler returned from `mysql_init()`.
- `rpl`: The replication stream structure. After a successful call, the `size` member indicates the event size, which is 0 for an EOF event. For a non-EOF event, `size` is greater than 0 and the `buffer` member points to the event contents.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

Example

See Section 27.7.17, “C API Binary Log Function Descriptions”.

27.7.17.3 mysql_binlog_open()

```
int mysql_binlog_open(MYSQL *mysql, MYSQL_RPL *rpl)
```

Description

Open a new replication event stream, to read a MySQL server binary log.

Arguments:

- `mysql`: The connection handler returned from `mysql_init()`.
- `rpl`: A `MYSQL_RPL` structure that has been initialized to indicate the replication event stream source. For a description of the structure members and how to initialize them, see Section 27.7.15, “C API Binary Log Data Structures”.

Return Values

Zero for success. Nonzero if an error occurred.

Errors

- `CR_FILE_NAME_TOO_LONG`

  The specified binary log file name was too long.

- `CR_OUT_OF_MEMORY`

  Out of memory.

Example

See Section 27.7.17, “C API Binary Log Function Descriptions”.

27.7.18 Common Questions and Problems When Using the C API

27.7.18.1 Why mysql_store_result() Sometimes Returns NULL After mysql_query() Returns Success

It is possible for `mysql_store_result()` to return `NULL` following a successful call to `mysql_query()`. When this happens, it means one of the following conditions occurred:

- There was a `malloc()` failure (for example, if the result set was too large).
- The data could not be read (an error occurred on the connection).

- The query returned no data (for example, it was an `INSERT`, `UPDATE`, or `DELETE`).

You can always check whether the statement should have produced a nonempty result by calling `mysql_field_count()`. If `mysql_field_count()` returns zero, the result is empty and the last query was a statement that does not return values (for example, an `INSERT` or a `DELETE`). If `mysql_field_count()` returns a nonzero value, the statement should have produced a nonempty result. See the description of the `mysql_field_count()` function for an example.

You can test for an error by calling `mysql_error()` or `mysql_errno()`.

27.7.18.2 What Results You Can Get from a Query

In addition to the result set returned by a query, you can also get the following information:

- `mysql_affected_rows()` returns the number of rows affected by the last query when doing an `INSERT`, `UPDATE`, or `DELETE`.

  For a fast re-create, use `TRUNCATE TABLE`.

- `mysql_num_rows()` returns the number of rows in a result set. With `mysql_store_result()`, `mysql_num_rows()` may be called as soon as `mysql_store_result()` returns. With `mysql_use_result()`, `mysql_num_rows()` may be called only after you have fetched all the rows with `mysql_fetch_row()`.

- `mysql_insert_id()` returns the ID generated by the last query that inserted a row into a table with an `AUTO_INCREMENT` index. See Section 27.7.7.38, “mysql_insert_id()”.

- Some queries (`LOAD DATA INFILE ...`, `INSERT INTO ... SELECT ...`, `UPDATE`) return additional information. The result is returned by `mysql_info()`. See the description for `mysql_info()` for the format of the string that it returns. `mysql_info()` returns a `NULL` pointer if there is no additional information.

27.7.18.3 How to Get the Unique ID for the Last Inserted Row

If you insert a record into a table that contains an `AUTO_INCREMENT` column, you can obtain the value stored into that column by calling the `mysql_insert_id()` function.

You can check from your C applications whether a value was stored in an `AUTO_INCREMENT` column by executing the following code (which assumes that you've checked that the statement succeeded). It determines whether the query was an `INSERT` with an `AUTO_INCREMENT` index:

```
if ((result = mysql_store_result(&mysql)) == 0 &&
    mysql_field_count(&mysql) == 0 &&
    mysql_insert_id(&mysql) != 0)
{
    used_id = mysql_insert_id(&mysql);
}
```

When a new `AUTO_INCREMENT` value has been generated, you can also obtain it by executing a `SELECT LAST_INSERT_ID()` statement with `mysql_query()` and retrieving the value from the result set returned by the statement.

When inserting multiple values, the last automatically incremented value is returned.

For `LAST_INSERT_ID()`, the most recently generated ID is maintained in the server on a per-connection basis. It is not changed by another client. It is not even changed if you update another `AUTO_INCREMENT` column with a nonmagic value (that is, a value that is not `NULL` and not `0`). Using `LAST_INSERT_ID()` and `AUTO_INCREMENT` columns simultaneously from multiple clients is perfectly valid. Each client will receive the last inserted ID for the last statement *that* client executed.

If you want to use the ID that was generated for one table and insert it into a second table, you can use SQL statements like this:

```
INSERT INTO foo (auto,text)
    VALUES(NULL,'text');         # generate ID by inserting NULL
INSERT INTO foo2 (id,text)
    VALUES(LAST_INSERT_ID(),'text');  # use ID in second table
```

`mysql_insert_id()` returns the value stored into an `AUTO_INCREMENT` column, whether that value is automatically generated by storing `NULL` or `0` or was specified as an explicit value. `LAST_INSERT_ID()` returns only automatically generated `AUTO_INCREMENT` values. If you store an explicit value other than `NULL` or `0`, it does not affect the value returned by `LAST_INSERT_ID()`.

For more information on obtaining the last ID in an `AUTO_INCREMENT` column:

- For information on `LAST_INSERT_ID()`, which can be used within an SQL statement, see Section 12.14, “Information Functions”.
- For information on `mysql_insert_id()`, the function you use from within the C API, see Section 27.7.7.38, “mysql_insert_id()”.
- For information on obtaining the auto-incremented value when using Connector/J, see Retrieving `AUTO_INCREMENT` Column Values through JDBC.
- For information on obtaining the auto-incremented value when using Connector/ODBC, see Obtaining Auto-Increment Values.

27.7.19 Controlling Automatic Reconnection Behavior

The MySQL client library can perform an automatic reconnection to the server if it finds that the connection is down when you attempt to send a statement to the server to be executed. If auto-reconnect is enabled, the library tries once to reconnect to the server and send the statement again.

Auto-reconnect is disabled by default.

If it is important for your application to know that the connection has been dropped (so that it can exit or take action to adjust for the loss of state information), be sure that auto-reconnect is disabled. To ensure this, call `mysql_options()` with the `MYSQL_OPT_RECONNECT` option:

```
bool reconnect = 0;
mysql_options(&mysql, MYSQL_OPT_RECONNECT, &reconnect);
```

If the connection has gone down, the effect of `mysql_ping()` depends on the auto-reconnect state. If auto-reconnect is enabled, `mysql_ping()` performs a reconnect. Otherwise, it returns an error.

Some client programs might provide the capability of controlling automatic reconnection. For example, `mysql` reconnects by default, but the `--skip-reconnect` option can be used to suppress this behavior.

If an automatic reconnection does occur (for example, as a result of calling `mysql_ping()`), there is no explicit indication of it. To check for reconnection, call `mysql_thread_id()` to get the original connection identifier before calling `mysql_ping()`, then call `mysql_thread_id()` again to see whether the identifier changed.

Automatic reconnection can be convenient because you need not implement your own reconnect code, but if a reconnection does occur, several aspects of the connection state are reset on the server side and your application will not be notified.

The connection-related state is affected as follows:

- Any active transactions are rolled back and autocommit mode is reset.
- All table locks are released.

- All `TEMPORARY` tables are closed (and dropped).
- Session system variables are reinitialized to the values of the corresponding global system variables, including system variables that are set implicitly by statements such as `SET NAMES`.
- User variable settings are lost.
- Prepared statements are released.
- `HANDLER` variables are closed.
- The value of `LAST_INSERT_ID()` is reset to 0.
- Locks acquired with `GET_LOCK()` are released.
- The association of the client with the Performance Schema `threads` table row that determines connection thread instrumentation is lost. If the client reconnects after a disconnect, the session is associated with a new row in the `threads` table and the thread monitoring state may be different. See Section 25.11.16.3, "The threads Table".

If the connection drops, it is possible that the session associated with the connection on the server side will still be running if the server has not yet detected that the client is no longer connected. In this case, any locks held by the original connection still belong to that session, so you may want to kill it by calling `mysql_kill()`.

27.7.20 C API Support for Multiple Statement Execution

By default, `mysql_query()` and `mysql_real_query()` interpret their statement string argument as a single statement to be executed, and you process the result according to whether the statement produces a result set (a set of rows, as for `SELECT`) or an affected-rows count (as for `INSERT`, `UPDATE`, and so forth).

MySQL also supports the execution of a string containing multiple statements separated by semicolon (`;`) characters. This capability is enabled by special options that are specified either when you connect to the server with `mysql_real_connect()` or after connecting by calling `mysql_set_server_option()`.

Executing a multiple-statement string can produce multiple result sets or row-count indicators. Processing these results involves a different approach than for the single-statement case: After handling the result from the first statement, it is necessary to check whether more results exist and process them in turn if so. To support multiple-result processing, the C API includes the `mysql_more_results()` and `mysql_next_result()` functions. These functions are used at the end of a loop that iterates as long as more results are available. *Failure to process the result this way may result in a dropped connection to the server.*

Multiple-result processing also is required if you execute `CALL` statements for stored procedures. Results from a stored procedure have these characteristics:

- Statements within the procedure may produce result sets (for example, if it executes `SELECT` statements). These result sets are returned in the order that they are produced as the procedure executes.

  In general, the caller cannot know how many result sets a procedure will return. Procedure execution may depend on loops or conditional statements that cause the execution path to differ from one call to the next. Therefore, you must be prepared to retrieve multiple results.

- The final result from the procedure is a status result that includes no result set. The status indicates whether the procedure succeeded or an error occurred.

The multiple statement and result capabilities can be used only with `mysql_query()` or `mysql_real_query()`. They cannot be used with the prepared statement interface. Prepared

statement handlers are defined to work only with strings that contain a single statement. See Section 27.7.8, “C API Prepared Statements”.

To enable multiple-statement execution and result processing, the following options may be used:

- The `mysql_real_connect()` function has a `flags` argument for which two option values are relevant:
  - `CLIENT_MULTI_RESULTS` enables the client program to process multiple results. This option *must* be enabled if you execute `CALL` statements for stored procedures that produce result sets. Otherwise, such procedures result in an error `Error 1312 (0A000): PROCEDURE proc_name can't return a result set in the given context`. `CLIENT_MULTI_RESULTS` is enabled by default.
  - `CLIENT_MULTI_STATEMENTS` enables `mysql_query()` and `mysql_real_query()` to execute statement strings containing multiple statements separated by semicolons. This option also enables `CLIENT_MULTI_RESULTS` implicitly, so a `flags` argument of `CLIENT_MULTI_STATEMENTS` to `mysql_real_connect()` is equivalent to an argument of `CLIENT_MULTI_STATEMENTS | CLIENT_MULTI_RESULTS`. That is, `CLIENT_MULTI_STATEMENTS` is sufficient to enable multiple-statement execution and all multiple-result processing.
- After the connection to the server has been established, you can use the `mysql_set_server_option()` function to enable or disable multiple-statement execution by passing it an argument of `MYSQL_OPTION_MULTI_STATEMENTS_ON` or `MYSQL_OPTION_MULTI_STATEMENTS_OFF`. Enabling multiple-statement execution with this function also enables processing of “simple” results for a multiple-statement string where each statement produces a single result, but is *not* sufficient to permit processing of stored procedures that produce result sets.

The following procedure outlines a suggested strategy for handling multiple statements:

1. Pass `CLIENT_MULTI_STATEMENTS` to `mysql_real_connect()`, to fully enable multiple-statement execution and multiple-result processing.
2. After calling `mysql_query()` or `mysql_real_query()` and verifying that it succeeds, enter a loop within which you process statement results.
3. For each iteration of the loop, handle the current statement result, retrieving either a result set or an affected-rows count. If an error occurs, exit the loop.
4. At the end of the loop, call `mysql_next_result()` to check whether another result exists and initiate retrieval for it if so. If no more results are available, exit the loop.

One possible implementation of the preceding strategy is shown following. The final part of the loop can be reduced to a simple test of whether `mysql_next_result()` returns nonzero. The code as written distinguishes between no more results and an error, which enables a message to be printed for the latter occurrence.

```
/* connect to server with the CLIENT_MULTI_STATEMENTS option */
if (mysql_real_connect (mysql, host_name, user_name, password,
    db_name, port_num, socket_name, CLIENT_MULTI_STATEMENTS) == NULL)
{
  printf("mysql_real_connect() failed\n");
  mysql_close(mysql);
  exit(1);
}

/* execute multiple statements */
status = mysql_query(mysql,
                     "DROP TABLE IF EXISTS test_table;\
                      CREATE TABLE test_table(id INT);\
                      INSERT INTO test_table VALUES(10);\
```

```
                      UPDATE test_table SET id=20 WHERE id=10;\
                      SELECT * FROM test_table;\
                      DROP TABLE test_table");
if (status)
{
  printf("Could not execute statement(s)");
  mysql_close(mysql);
  exit(0);
}

/* process each statement result */
do {
  /* did current statement return data? */
  result = mysql_store_result(mysql);
  if (result)
  {
    /* yes; process rows and free the result set */
    process_result_set(mysql, result);
    mysql_free_result(result);
  }
  else          /* no result set or error */
  {
    if (mysql_field_count(mysql) == 0)
    {
      printf("%lld rows affected\n",
             mysql_affected_rows(mysql));
    }
    else  /* some error occurred */
    {
      printf("Could not retrieve result set\n");
      break;
    }
  }
  /* more results? -1 = no, >0 = error, 0 = yes (keep looping) */
  if ((status = mysql_next_result(mysql)) > 0)
    printf("Could not execute statement\n");
} while (status == 0);

mysql_close(mysql);
```

27.7.21 C API Prepared Statement Problems

Here follows a list of the currently known problems with prepared statements:

- `TIME`, `TIMESTAMP`, and `DATETIME` do not support parts of seconds (for example, from `DATE_FORMAT()`).
- When converting an integer to string, `ZEROFILL` is honored with prepared statements in some cases where the MySQL server does not print the leading zeros. (For example, with `MIN(`*`number-with-zerofill`*`)`).
- When converting a floating-point number to a string in the client, the rightmost digits of the converted value may differ slightly from those of the original value.
- Prepared statements use the query cache under the conditions described in Section 8.10.3.1, “How the Query Cache Operates”.
- Prepared statements do not support multi-statements (that is, multiple statements within a single string separated by `;` characters).
- The capabilities of prepared `CALL` statements are described in Section 27.7.23, “C API Support for Prepared CALL Statements”.

27.7.22 C API Prepared Statement Handling of Date and Time Values

The binary (prepared statement) protocol enables you to send and receive date and time values (`DATE`, `TIME`, `DATETIME`, and `TIMESTAMP`), using the `MYSQL_TIME` structure. The members of this structure are described in Section 27.7.9, “C API Prepared Statement Data Structures”.

To send temporal data values, create a prepared statement using `mysql_stmt_prepare()`. Then, before calling `mysql_stmt_execute()` to execute the statement, use the following procedure to set up each temporal parameter:

1. In the `MYSQL_BIND` structure associated with the data value, set the `buffer_type` member to the type that indicates what kind of temporal value you're sending. For `DATE`, `TIME`, `DATETIME`, or `TIMESTAMP` values, set `buffer_type` to `MYSQL_TYPE_DATE`, `MYSQL_TYPE_TIME`, `MYSQL_TYPE_DATETIME`, or `MYSQL_TYPE_TIMESTAMP`, respectively.

2. Set the `buffer` member of the `MYSQL_BIND` structure to the address of the `MYSQL_TIME` structure in which you pass the temporal value.

3. Fill in the members of the `MYSQL_TIME` structure that are appropriate for the type of temporal value to pass.

Use `mysql_stmt_bind_param()` to bind the parameter data to the statement. Then you can call `mysql_stmt_execute()`.

To retrieve temporal values, the procedure is similar, except that you set the `buffer_type` member to the type of value you expect to receive, and the `buffer` member to the address of a `MYSQL_TIME` structure into which the returned value should be placed. Use `mysql_stmt_bind_result()` to bind the buffers to the statement after calling `mysql_stmt_execute()` and before fetching the results.

Here is a simple example that inserts `DATE`, `TIME`, and `TIMESTAMP` data. The `mysql` variable is assumed to be a valid connection handler.

```
MYSQL_TIME  ts;
MYSQL_BIND  bind[3];
MYSQL_STMT  *stmt;

strmov(query, "INSERT INTO test_table(date_field, time_field, \
                            timestamp_field) VALUES(?,?,?");

stmt = mysql_stmt_init(mysql);
if (!stmt)
{
  fprintf(stderr, " mysql_stmt_init(), out of memory\n");
  exit(0);
}
if (mysql_stmt_prepare(mysql, query, strlen(query)))
{
  fprintf(stderr, "\n mysql_stmt_prepare(), INSERT failed");
  fprintf(stderr, "\n %s", mysql_stmt_error(stmt));
  exit(0);
}

/* set up input buffers for all 3 parameters */
bind[0].buffer_type= MYSQL_TYPE_DATE;
bind[0].buffer= (char *)&ts;
bind[0].is_null= 0;
bind[0].length= 0;
...
bind[1]= bind[2]= bind[0];
...

mysql_stmt_bind_param(stmt, bind);

/* supply the data to be sent in the ts structure */
ts.year= 2002;
ts.month= 02;
ts.day= 03;

ts.hour= 10;
ts.minute= 45;
ts.second= 20;
```

```
mysql_stmt_execute(stmt);
..
```

27.7.23 C API Support for Prepared CALL Statements

This section describes prepared-statement support in the C API for stored procedures executed using `CALL` statements:

Stored procedures executed using prepared `CALL` statements can be used in the following ways:

- A stored procedure can produce any number of result sets. The number of columns and the data types of the columns need not be the same for all result sets.
- The final values of `OUT` and `INOUT` parameters are available to the calling application after the procedure returns. These parameters are returned as an extra single-row result set following any result sets produced by the procedure itself. The row contains the values of the `OUT` and `INOUT` parameters in the order in which they are declared in the procedure parameter list.

The following discussion shows how to use these capabilities through the C API for prepared statements. To use prepared `CALL` statements through the `PREPARE` and `EXECUTE` statements, see Section 13.2.1, “CALL Syntax”.

If an application might be compiled or executed in a context where a version of MySQL older than 5.5.3 is used, prepared `CALL` capabilities for multiple result sets and `OUT` or `INOUT` parameters might not be available:

- For the client side, the application will not compile unless the libraries are from MySQL 5.5.3 or higher (the API function and symbols introduced in that version will not be present).
- To verify at runtime that the server is recent enough, a client can use this test:

```
if (mysql_get_server_version(mysql) < 50503)
{
  fprintf(stderr,
          "Server does not support required CALL capabilities\n");
  mysql_close(mysql);
  exit (1);
}
```

An application that executes a prepared `CALL` statement should use a loop that fetches a result and then invokes `mysql_stmt_next_result()` to determine whether there are more results. The results consist of any result sets produced by the stored procedure followed by a final status value that indicates whether the procedure terminated successfully.

If the procedure has `OUT` or `INOUT` parameters, the result set preceding the final status value contains their values. To determine whether a result set contains parameter values, test whether the `SERVER_PS_OUT_PARAMS` bit is set in the `server_status` member of the `MYSQL` connection handler:

```
mysql->server_status & SERVER_PS_OUT_PARAMS
```

The following example uses a prepared `CALL` statement to execute a stored procedure that produces multiple result sets and that provides parameter values back to the caller by means of `OUT` and `INOUT` parameters. The procedure takes parameters of all three types (`IN`, `OUT`, `INOUT`), displays their initial values, assigns new values, displays the updated values, and returns. The expected return information from the procedure therefore consists of multiple result sets and a final status:

- One result set from a `SELECT` that displays the initial parameter values: `10`, `NULL`, `30`. (The `OUT` parameter is assigned a value by the caller, but this assignment is expected to be ineffective: `OUT` parameters are seen as `NULL` within a procedure until assigned a value within the procedure.)
- One result set from a `SELECT` that displays the modified parameter values: `100`, `200`, `300`.

- One result set containing the final `OUT` and `INOUT` parameter values: `200`, `300`.
- A final status packet.

The code to execute the procedure:

```
MYSQL_STMT *stmt;
MYSQL_BIND ps_params[3];  /* input parameter buffers */
int        int_data[3];   /* input/output values */
bool       is_null[3];    /* output value nullability */
int        status;

/* set up stored procedure */
status = mysql_query(mysql, "DROP PROCEDURE IF EXISTS p1");
test_error(mysql, status);

status = mysql_query(mysql,
  "CREATE PROCEDURE p1("
  "  IN p_in INT, "
  "  OUT p_out INT, "
  "  INOUT p_inout INT) "
  "BEGIN "
  "  SELECT p_in, p_out, p_inout; "
  "  SET p_in = 100, p_out = 200, p_inout = 300; "
  "  SELECT p_in, p_out, p_inout; "
  "END");
test_error(mysql, status);

/* initialize and prepare CALL statement with parameter placeholders */
stmt = mysql_stmt_init(mysql);
if (!stmt)
{
  printf("Could not initialize statement\n");
  exit(1);
}
status = mysql_stmt_prepare(stmt, "CALL p1(?, ?, ?)", 16);
test_stmt_error(stmt, status);

/* initialize parameters: p_in, p_out, p_inout (all INT) */
memset(ps_params, 0, sizeof (ps_params));

ps_params[0].buffer_type = MYSQL_TYPE_LONG;
ps_params[0].buffer = (char *) &int_data[0];
ps_params[0].length = 0;
ps_params[0].is_null = 0;

ps_params[1].buffer_type = MYSQL_TYPE_LONG;
ps_params[1].buffer = (char *) &int_data[1];
ps_params[1].length = 0;
ps_params[1].is_null = 0;

ps_params[2].buffer_type = MYSQL_TYPE_LONG;
ps_params[2].buffer = (char *) &int_data[2];
ps_params[2].length = 0;
ps_params[2].is_null = 0;

/* bind parameters */
status = mysql_stmt_bind_param(stmt, ps_params);
test_stmt_error(stmt, status);

/* assign values to parameters and execute statement */
int_data[0]= 10;  /* p_in */
int_data[1]= 20;  /* p_out */
int_data[2]= 30;  /* p_inout */

status = mysql_stmt_execute(stmt);
test_stmt_error(stmt, status);

/* process results until there are no more */
do {
  int i;
```

```
int num_fields;       /* number of columns in result */
MYSQL_FIELD *fields;  /* for result set metadata */
MYSQL_BIND *rs_bind;  /* for output buffers */

/* the column count is > 0 if there is a result set */
/* 0 if the result is only the final status packet */
num_fields = mysql_stmt_field_count(stmt);

if (num_fields > 0)
{
  /* there is a result set to fetch */
  printf("Number of columns in result: %d\n", (int) num_fields);

  /* what kind of result set is this? */
  printf("Data: ");
  if(mysql->server_status & SERVER_PS_OUT_PARAMS)
    printf("this result set contains OUT/INOUT parameters\n");
  else
    printf("this result set is produced by the procedure\n");

  MYSQL_RES *rs_metadata = mysql_stmt_result_metadata(stmt);
  test_stmt_error(stmt, rs_metadata == NULL);

  fields = mysql_fetch_fields(rs_metadata);

  rs_bind = (MYSQL_BIND *) malloc(sizeof (MYSQL_BIND) * num_fields);
  if (!rs_bind)
  {
    printf("Cannot allocate output buffers\n");
    exit(1);
  }
  memset(rs_bind, 0, sizeof (MYSQL_BIND) * num_fields);

  /* set up and bind result set output buffers */
  for (i = 0; i < num_fields; ++i)
  {
    rs_bind[i].buffer_type = fields[i].type;
    rs_bind[i].is_null = &is_null[i];

    switch (fields[i].type)
    {
      case MYSQL_TYPE_LONG:
        rs_bind[i].buffer = (char *) &(int_data[i]);
        rs_bind[i].buffer_length = sizeof (int_data);
        break;

      default:
        fprintf(stderr, "ERROR: unexpected type: %d.\n", fields[i].type);
        exit(1);
    }
  }

  status = mysql_stmt_bind_result(stmt, rs_bind);
  test_stmt_error(stmt, status);

  /* fetch and display result set rows */
  while (1)
  {
    status = mysql_stmt_fetch(stmt);

    if (status == 1 || status == MYSQL_NO_DATA)
      break;

    for (i = 0; i < num_fields; ++i)
    {
      switch (rs_bind[i].buffer_type)
      {
        case MYSQL_TYPE_LONG:
          if (*rs_bind[i].is_null)
            printf(" val[%d] = NULL;", i);
          else
            printf(" val[%d] = %ld;",
```

```
                    i, (long) *((int *) rs_bind[i].buffer));
            break;

          default:
            printf("  unexpected type (%d)\n",
              rs_bind[i].buffer_type);
        }
      }
      printf("\n");
    }

    mysql_free_result(rs_metadata); /* free metadata */
    free(rs_bind);                  /* free output buffers */
  }
  else
  {
    /* no columns = final status packet */
    printf("End of procedure output\n");
  }

  /* more results? -1 = no, >0 = error, 0 = yes (keep looking) */
  status = mysql_stmt_next_result(stmt);
  if (status > 0)
    test_stmt_error(stmt, status);
} while (status == 0);

mysql_stmt_close(stmt);
```

Execution of the procedure should produce the following output:

```
Number of columns in result: 3
Data: this result set is produced by the procedure
 val[0] = 10; val[1] = NULL; val[2] = 30;
Number of columns in result: 3
Data: this result set is produced by the procedure
 val[0] = 100; val[1] = 200; val[2] = 300;
Number of columns in result: 2
Data: this result set contains OUT/INOUT parameters
 val[0] = 200; val[1] = 300;
End of procedure output
```

The code uses two utility routines, `test_error()` and `test_stmt_error()`, to check for errors and terminate after printing diagnostic information if an error occurred:

```
static void test_error(MYSQL *mysql, int status)
{
  if (status)
  {
    fprintf(stderr, "Error: %s (errno: %d)\n",
            mysql_error(mysql), mysql_errno(mysql));
    exit(1);
  }
}

static void test_stmt_error(MYSQL_STMT *stmt, int status)
{
  if (status)
  {
    fprintf(stderr, "Error: %s (errno: %d)\n",
            mysql_stmt_error(stmt), mysql_stmt_errno(stmt));
    exit(1);
  }
}
```

27.8 MySQL PHP API

The MySQL PHP API manual is now published in standalone form, not as part of the MySQL Reference Manual. See MySQL and PHP.

27.9 MySQL Perl API

The Perl `DBI` module provides a generic interface for database access. You can write a DBI script that works with many different database engines without change. To use DBI with MySQL, install the following:

1. The `DBI` module.
2. The `DBD::mysql` module. This is the DataBase Driver (DBD) module for Perl.
3. Optionally, the DBD module for any other type of database server you want to access.

Perl DBI is the recommended Perl interface. It replaces an older interface called `mysqlperl`, which should be considered obsolete.

These sections contain information about using Perl with MySQL and writing MySQL applications in Perl:

- For installation instructions for Perl DBI support, see Section 2.11, “Perl Installation Notes”.
- For an example of reading options from option files, see Section 5.7.4, “Using Client Programs in a Multiple-Server Environment”.
- For secure coding tips, see Section 6.1.1, “Security Guidelines”.
- For debugging tips, see Section 28.5.1.4, “Debugging mysqld under gdb”.
- For some Perl-specific environment variables, see Section 4.9, “MySQL Program Environment Variables”.
- For considerations for running on OS X, see Section 2.4, “Installing MySQL on OS X”.
- For ways to quote string literals, see Section 9.1.1, “String Literals”.

DBI information is available at the command line, online, or in printed form:

- Once you have the `DBI` and `DBD::mysql` modules installed, you can get information about them at the command line with the `perldoc` command:

  ```
  shell> perldoc DBI
  shell> perldoc DBI::FAQ
  shell> perldoc DBD::mysql
  ```

  You can also use `pod2man`, `pod2html`, and so on to translate this information into other formats.

- For online information about Perl DBI, visit the DBI Web site, http://dbi.perl.org/. That site hosts a general DBI mailing list. Oracle Corporation hosts a list specifically about `DBD::mysql`; see Section 1.6.2, “MySQL Mailing Lists”.

- For printed information, the official DBI book is *Programming the Perl DBI* (Alligator Descartes and Tim Bunce, O'Reilly & Associates, 2000). Information about the book is available at the DBI Web site, http://dbi.perl.org/.

  For information that focuses specifically on using DBI with MySQL, see *MySQL and Perl for the Web* (Paul DuBois, New Riders, 2001). This book's Web site is http://www.kitebird.com/mysql-perl/.

27.10 MySQL Python API

`MySQLdb` is a third-party driver that provides MySQL support for Python, compliant with the Python DB API version 2.0. It can be found at http://sourceforge.net/projects/mysql-python/.

The new MySQL Connector/Python component provides an interface to the same Python API, and is built into the MySQL Server and supported by Oracle. See MySQL Connector/Python Developer Guide for details on the Connector, as well as coding guidelines for Python applications and sample Python code.

27.11 MySQL Ruby APIs

Two APIs are available for Ruby programmers developing MySQL applications:

- The MySQL/Ruby API is based on the `libmysqlclient` API library. For information on installing and using the MySQL/Ruby API, see Section 27.11.1, “The MySQL/Ruby API”.
- The Ruby/MySQL API is written to use the native MySQL network protocol (a native driver). For information on installing and using the Ruby/MySQL API, see Section 27.11.2, “The Ruby/MySQL API”.

For background and syntax information about the Ruby language, see Ruby Programming Language.

27.11.1 The MySQL/Ruby API

The MySQL/Ruby module provides access to MySQL databases using Ruby through `libmysqlclient`.

For information on installing the module, and the functions exposed, see MySQL/Ruby.

27.11.2 The Ruby/MySQL API

The Ruby/MySQL module provides access to MySQL databases using Ruby through a native driver interface using the MySQL network protocol.

For information on installing the module, and the functions exposed, see Ruby/MySQL.

27.12 MySQL Tcl API

`MySQLtcl` is a simple API for accessing a MySQL database server from the Tcl programming language. It can be found at http://www.xdobry.de/mysqltcl/.

27.13 MySQL Eiffel Wrapper

Eiffel MySQL is an interface to the MySQL database server using the Eiffel programming language, written by Michael Ravits. It can be found at http://efsa.sourceforge.net/archive/ravits/mysql.htm.

Chapter 28 Extending MySQL

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28.1 MySQL Internals

This chapter describes a lot of things that you need to know when working on the MySQL code. To track or contribute to MySQL development, follow the instructions in Section 2.8.3, “Installing MySQL Using a Development Source Tree”. If you are interested in MySQL internals, you should also subscribe to our `internals` mailing list. This list has relatively low traffic. For details on how to subscribe, please see Section 1.6.2, “MySQL Mailing Lists”. Many MySQL developers at Oracle Corporation are on the `internals` list and we help other people who are working on the MySQL code. Feel free to use this list both to ask questions about the code and to send patches that you would like to contribute to the MySQL project!

Note

The MySQL source code contains internal documentation written using Doxygen. This documentation is useful for understanding how MySQL works from a developer perspective. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, “Generating MySQL Doxygen Documentation Content”.

28.1.1 MySQL Threads

The MySQL server creates the following threads:

- Connection manager threads handle client connection requests on the network interfaces that the server listens to. On all platforms, one manager thread handles TCP/IP connection requests. On Unix, this manager thread also handles Unix socket file connection requests. On Windows, a manager thread handles shared-memory connection requests, and another handles named-pipe connection requests. The server does not create threads to handle interfaces that it does not listen to. For example, a Windows server that does not have support for named-pipe connections enabled does not create a thread to handle them.

- Connection manager threads associate each client connection with a thread dedicated to it that handles authentication and request processing for that connection. Manager threads create a new thread when necessary but try to avoid doing so by consulting the thread cache first to see whether it contains a thread that can be used for the connection. When a connection ends, its thread is returned to the thread cache if the cache is not full.

  For information about tuning the parameters that control thread resources, see Section 8.12.4.1, "How MySQL Uses Threads for Client Connections".

- On a master replication server, connections from slave servers are handled like client connections: There is one thread per connected slave.

- On a slave replication server, an I/O thread is started to connect to the master server and read updates from it. An SQL thread is started to apply updates read from the master. These two threads run independently and can be started and stopped independently.

- A signal thread handles all signals. This thread also normally handles alarms and calls `process_alarm()` to force timeouts on connections that have been idle too long.

- If `InnoDB` is used, there will be additional read and write threads by default. The number of these are controlled by the `innodb_read_io_threads` and `innodb_write_io_threads` parameters. See Section 15.13, "InnoDB Startup Options and System Variables".

- If the server is started with the `--flush_time=val` option, a dedicated thread is created to flush all tables every *`val`* seconds.

- If the event scheduler is active, there is one thread for the scheduler, and a thread for each event currently running. See Section 23.4.1, "Event Scheduler Overview".

`mysqladmin processlist` only shows the connection, replication, and event threads.

28.1.2 The MySQL Test Suite

The test system that is included in Unix source and binary distributions makes it possible for users and developers to perform regression tests on the MySQL code. These tests can be run on Unix.

You can also write your own test cases. For information about the MySQL Test Framework, including system requirements, see the manual available at http://dev.mysql.com/doc/mysqltest/2.0/en/.

The current set of test cases does not test everything in MySQL, but it should catch most obvious bugs in the SQL processing code, operating system or library issues, and is quite thorough in testing replication. Our goal is to have the tests cover 100% of the code. We welcome contributions to our test suite. You may especially want to contribute tests that examine the functionality critical to your system because this ensures that all future MySQL releases work well with your applications.

The test system consists of a test language interpreter (`mysqltest`), a Perl script to run all tests (`mysql-test-run.pl`), the actual test cases written in a special test language, and their expected results. To run the test suite on your system after a build, type `make test` from the source root directory, or change location to the `mysql-test` directory and type `./mysql-test-run.pl`. If you have installed a binary distribution, change location to the `mysql-test` directory under the installation root directory (for example, `/usr/local/mysql/mysql-test`), and run `./mysql-test-run.pl`. All tests should succeed. If any do not, feel free to try to find out why and report the problem if it indicates a bug in MySQL. See Section 1.7, "How to Report Bugs or Problems".

If one test fails, you should run `mysql-test-run.pl` with the `--force` option to check whether any other tests fail.

If you have a copy of `mysqld` running on the machine where you want to run the test suite, you do not have to stop it, as long as it is not using ports `9306` or `9307`. If either of those ports is taken, you should set the `MTR_BUILD_THREAD` environment variable to an appropriate value, and the test suite will use a different set of ports for master, slave, and NDB). For example:

```
shell> export MTR_BUILD_THREAD=31
shell> ./mysql-test-run.pl [options] [test_name]
```

In the `mysql-test` directory, you can run an individual test case with `./mysql-test-run.pl test_name`.

If you have a question about the test suite, or have a test case to contribute, send an email message to the MySQL `internals` mailing list. See Section 1.6.2, “MySQL Mailing Lists”.

28.2 The MySQL Plugin API

MySQL supports a plugin API that enables creation of server components. Plugins can be loaded at server startup, or loaded and unloaded at runtime without restarting the server. The API is generic and does not specify what plugins can do. The components supported by this interface include, but are not limited to, storage engines, full-text parser plugins, and server extensions.

For example, full-text parser plugins can be used to replace or augment the built-in full-text parser. A plugin can parse text into words using rules that differ from those used by the built-in parser. This can be useful if you need to parse text with characteristics different from those expected by the built-in parser.

The plugin interface is more general than the older user-defined function (UDF) interface.

The plugin interface uses the `plugin` table in the `mysql` database to record information about plugins that have been installed permanently with the `INSTALL PLUGIN` statement. This table is created as part of the MySQL installation process. Plugins can also be installed for a single server invocation with the `--plugin-load` option. Plugins installed this way are not recorded in the `plugin` table. See Section 5.6.2, “Installing and Uninstalling Plugins”.

MySQL supports an API for client plugins in addition to that for server plugins. This is used, for example, by authentication plugins where a server-side plugin and a client-side plugin cooperate to enable clients to connect to the server through a variety of authentication methods.

Note

The MySQL source code contains internal documentation written using Doxygen. This documentation is useful for understanding how MySQL works from a developer perspective. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, “Generating MySQL Doxygen Documentation Content”.

Additional Resources

The book *MySQL 5.1 Plugin Development* by Sergei Golubchik and Andrew Hutchings provides a wealth of detail about the plugin API. Despite the fact that the book's title refers to MySQL Server 5.1, most of the information in it applies to later versions as well.

28.2.1 Types of Plugins

The plugin API enables creation of plugins that implement several capabilities:

- Storage engines
- Full-text parsers
- Daemons
- `INFORMATION_SCHEMA` tables

- Semisynchronous replication
- Auditing
- Authentication
- Password validation and strength checking
- Protocol tracing
- Query rewriting
- Secure keyring storage and retrieval

The following sections provide an overview of these plugin types.

Storage Engine Plugins

The pluggable storage engine architecture used by MySQL Server enables storage engines to be written as plugins and loaded into and unloaded from a running server. For a description of this architecture, see Section 16.11, “Overview of MySQL Storage Engine Architecture”.

For information on how to use the plugin API to write storage engines, see MySQL Internals: Writing a Custom Storage Engine.

Full-Text Parser Plugins

MySQL has a built-in parser that it uses by default for full-text operations (parsing text to be indexed, or parsing a query string to determine the terms to be used for a search). The built-in full-text parser is supported with `InnoDB` and `MyISAM` tables.

MySQL also has a character-based ngram full-text parser that supports Chinese, Japanese, and Korean (CJK), and a word-based MeCab parser plugin that supports Japanese, for use with `InnoDB` and `MyISAM` tables.

For full-text processing, “parsing” means extracting words (or “tokens”, in the case of an n-gram character-based parser) from text or a query string based on rules that define which character sequences make up a word and where word boundaries lie.

When parsing for indexing purposes, the parser passes each word to the server, which adds it to a full-text index. When parsing a query string, the parser passes each word to the server, which accumulates the words for use in a search.

The parsing properties of the built-in full-text parser are described in Section 12.9, “Full-Text Search Functions”. These properties include rules for determining how to extract words from text. The parser is influenced by certain system variables that cause words shorter or longer to be excluded, and by the stopword list that identifies common words to be ignored. For more information, see Section 12.9.4, “Full-Text Stopwords”, and Section 12.9.6, “Fine-Tuning MySQL Full-Text Search”.

The plugin API enables you to use a full-text parser other than the default built-in full-text parser. For example, if you are working with Japanese, you may choose to use the MeCab full-text parser. The plugin API also enables you to provide a full-text parser of your own so that you have control over the basic duties of a parser. A parser plugin can operate in either of two roles:

- The plugin can replace the built-in parser. In this role, the plugin reads the input to be parsed, splits it up into words, and passes the words to the server (either for indexing or for token accumulation). The ngram and MeCab parsers operate as replacements for the built-in full-text parser.

  You may choose to provide your own full-text parser if you need to use different rules from those of the built-in parser for determining how to split up input into words. For example, the built-in parser considers the text “case-sensitive” to consist of two words “case” and “sensitive,” whereas an application might need to treat the text as a single word.

- The plugin can act in conjunction with the built-in parser by serving as a front end for it. In this role, the plugin extracts text from the input and passes the text to the parser, which splits up the text into words using its normal parsing rules. This parsing is affected by the `innodb_ft_xxx` or `ft_xxx` system variables and the stopword list.

  One reason to use a parser this way is that you need to index content such as PDF documents, XML documents, or `.doc` files. The built-in parser is not intended for those types of input but a plugin can pull out the text from these input sources and pass it to the built-in parser.

It is also possible for a parser plugin to operate in both roles. That is, it could extract text from noncleartext input (the front end role), and also parse the text into words (thus replacing the built-in parser).

A full-text plugin is associated with full-text indexes on a per-index basis. That is, when you install a parser plugin initially, that does not cause it to be used for any full-text operations. It simply becomes available. For example, a full-text parser plugin becomes available to be named in a `WITH PARSER` clause when creating individual `FULLTEXT` indexes. To create such an index at table-creation time, do this:

```
CREATE TABLE t
(
  doc CHAR(255),
  FULLTEXT INDEX (doc) WITH PARSER parser_name
) ENGINE=InnoDB;
```

Or you can add the index after the table has been created:

```
ALTER TABLE t ADD FULLTEXT INDEX (doc) WITH PARSER parser_name;
```

The only SQL change for associating the parser with the index is the `WITH PARSER` clause. Searches are specified as before, with no changes needed for queries.

When you associate a parser plugin with a `FULLTEXT` index, the plugin is required for using the index. If the parser plugin is dropped, any index associated with it becomes unusable. Any attempt to use a table for which a plugin is not available results in an error, although `DROP TABLE` is still possible.

For more information about full-text plugins, see Section 28.2.4.4, “Writing Full-Text Parser Plugins”. MySQL 8.0 supports full-text plugins with `MyISAM` and `InnoDB`.

Daemon Plugins

A daemon plugin is a simple type of plugin used for code that should be run by the server but that does not communicate with it. MySQL distributions include an example daemon plugin that writes periodic heartbeat messages to a file.

For more information about daemon plugins, see Section 28.2.4.5, “Writing Daemon Plugins”.

INFORMATION_SCHEMA Plugins

`INFORMATION_SCHEMA` plugins enable the creation of tables containing server metadata that are exposed to users through the `INFORMATION_SCHEMA` database. For example, `InnoDB` uses `INFORMATION_SCHEMA` plugins to provide tables that contain information about current transactions and locks.

For more information about `INFORMATION_SCHEMA` plugins, see Section 28.2.4.6, “Writing INFORMATION_SCHEMA Plugins”.

Semisynchronous Replication Plugins

MySQL replication is asynchronous by default. With semisynchronous replication, a commit performed on the master side blocks before returning to the session that performed the transaction until

at least one slave acknowledges that it has received and logged the events for the transaction. Semisynchronous replication is implemented through complementary master and client plugins. See Section 18.3.10, “Semisynchronous Replication”.

For more information about semisynchronous replication plugins, see Section 28.2.4.7, “Writing Semisynchronous Replication Plugins”.

Audit Plugins

The MySQL server provides a pluggable audit interface that enables information about server operations to be reported to interested parties. Audit notification occurs for these operations (although the interface is general and the server could be modified to report others):

- Write a message to the general query log (if the log is enabled)
- Write a message to the error log
- Send a query result to a client

Audit plugins may register with the audit interface to receive notification about server operations. When an auditable event occurs within the server, the server determines whether notification is needed. For each registered audit plugin, the server checks the event against those event classes in which the plugin is interested and passes the event to the plugin if there is a match.

This interface enables audit plugins to receive notifications only about operations in event classes they consider significant and to ignore others. The interface provides for categorization of operations into event classes and further division into event subclasses within each class.

When an audit plugin is notified of an auditable event, it receives a pointer to the current THD structure and a pointer to a structure that contains information about the event. The plugin can examine the event and perform whatever auditing actions are appropriate. For example, the plugin can see what statement produced a result set or was logged, the number of rows in a result, who the current user was for an operation, or the error code for failed operations.

For more information about audit plugins, see Section 28.2.4.8, “Writing Audit Plugins”.

Authentication Plugins

MySQL supports pluggable authentication. Authentication plugins exist on both the server and client sides. Plugins on the server side implement authentication methods for use by clients when they connect to the server. A plugin on the client side communicates with a server-side plugin to provide the authentication information that it requires. A client-side plugin may interact with the user, performing tasks such as soliciting a password or other authentication credentials to be sent to the server. See Section 6.3.10, “Pluggable Authentication”.

Pluggable authentication also enables proxy user capability, in which one user takes the identity of another user. A server-side authentication plugin can return to the server the name of the user whose identity the connecting user should have. See Section 6.3.11, “Proxy Users”.

For more information about authentication plugins, see Section 28.2.4.9, “Writing Authentication Plugins”.

Password-Validation Plugins

The MySQL server provides an interface for writing plugins that test passwords. Such a plugin implements two capabilities:

- Rejection of too-weak passwords in statements that assign passwords (such as `CREATE USER`, `GRANT`, and `SET PASSWORD` statements), and passwords given as arguments to the `PASSWORD()` function.

- Assessing the strength of potential passwords for the `VALIDATE_PASSWORD_STRENGTH()` SQL function.

For information about writing this type of plugin, see Section 28.2.4.10, “Writing Password-Validation Plugins”.

Protocol Trace Plugins

MySQL supports the use of protocol trace plugins: client-side plugins that implement tracing of communication between a client and the server that takes place using the client/server protocol.

For more information about protocol trace plugins, see Section 28.2.4.11, “Writing Protocol Trace Plugins”.

Query Rewrite Plugins

MySQL Server supports query rewrite plugins that can examine and possibly modify statements received by the server before the server executes them. A query rewrite plugin takes statements either before or after the server has parsed them.

A preparse query rewrite plugin has these characteristics:

- The plugin enables rewriting of SQL statements arriving at the server before the server processes them.
- The plugin receives a statement string and may return a different string.

A postparse query rewrite plugin has these characteristics:

- The plugin enables statement rewriting based on parse trees.
- The server parses each statement and passes its parse tree to the plugin, which may traverse the tree. The plugin can return the original tree to the server for further processing, or construct a different tree and return that instead.
- The plugin can use the `mysql_parser` plugin service for these purposes:
  - To activate statement digest calculation and obtain the normalized version of statements independent of whether the Performance Schema produces digests.
  - To traverse parse trees.
  - To parse statements. This is useful if the plugin constructs a new statement string from the parse tree. The plugin can have the server parse the string to produce a new tree, then return that tree as the representation of the rewritten statement.

For more information about plugin services, see Section 28.3, “MySQL Services for Components and Plugins”.

Preparse and postparse query rewrite plugins share these characteristics:

- If a query rewrite plugin is installed, the `--log-raw` option affects statement logging as follows:
  - Without `--log-raw`, the server logs the statement returned by the query rewrite plugin. This may differ from the statement as received.
  - With `--log-raw`, the server logs the original statement as received.
- If a plugin rewrites a statement, the server decides whether to write it to the binary log (and thus to any replication slaves) based on the rewritten statement, not the original statement. If a plugin rewrites only `SELECT` statements to `SELECT` statements, there is no impact on binary logging because the server does not write `SELECT` statements to the binary log.

- If a plugin rewrites a statement, the server produces a `Note` message that the client can view using `SHOW WARNINGS`. Messages have this format, where *`stmt_in`* is the original statement and *`stmt_out`* is the rewritten statement:

```
Query 'stmt_in' rewritten to 'stmt_out' by a query rewrite plugin
```

MySQL distributions include a postparse query rewrite plugin named `Rewriter`. This plugin is rule based. You can add rows to its rules table to cause `SELECT` statement rewriting. For more information, see Section 5.6.4, “The Rewriter Query Rewrite Plugin”.

Query rewrite plugins use the same API as audit plugins. For more information about audit plugins, see Section 28.2.4.8, “Writing Audit Plugins”.

Keyring Plugins

MySQL Server supports keyring plugins that enable internal server components and plugins to securely store sensitive information for later retrieval.

All MySQL distributions include a keyring plugin named `keyring_file`. See Section 6.5.4, “The MySQL Keyring”.

For more information about keyring plugins, see Section 28.2.4.12, “Writing Keyring Plugins”.

28.2.2 Plugin API Characteristics

The server plugin API has these characteristics:

- All plugins have several things in common.

  Each plugin has a name that it can be referred to in SQL statements, as well as other metadata such as an author and a description that provide other information. This information can be examined in the `INFORMATION_SCHEMA.PLUGINS` table or using the `SHOW PLUGINS` statement.

- The plugin framework is extendable to accommodate different kinds of plugins.

  Although some aspects of the plugin API are common to all types of plugins, the API also permits type-specific interface elements so that different types of plugins can be created. A plugin with one purpose can have an interface most appropriate to its own requirements and not the requirements of some other plugin type.

  Interfaces for several types of plugins exist, such as storage engines, full-text parser, and `INFORMATION_SCHEMA` tables. Others can be added.

- Plugins can expose information to users.

  A plugin can implement system and status variables that are available through the `SHOW VARIABLES` and `SHOW STATUS` statements.

- The plugin API includes versioning information.

  The version information included in the plugin API enables a plugin library and each plugin that it contains to be self-identifying with respect to the API version that was used to build the library. If the API changes over time, the version numbers will change, but a server can examine a given plugin library's version information to determine whether it supports the plugins in the library.

  There are two types of version numbers. The first is the version for the general plugin framework itself. Each plugin library includes this kind of version number. The second type of version applies to individual plugins. Each specific type of plugin has a version for its interface, so each plugin in a library has a type-specific version number. For example, a library containing a full-text parser plugin has a general plugin API version number, and the plugin has a version number specific to the full-text plugin interface.

- The plugin API implements security restrictions.

  A plugin library must be installed in a specific dedicated directory for which the location is controlled by the server and cannot be changed at runtime. Also, the library must contain specific symbols that identify it as a plugin library. The server will not load something as a plugin if it was not built as a plugin.

- Plugins have access to server services.

  The services interface exposes server functionality that plugins can access using ordinary function calls. For details, see Section 28.3, “MySQL Services for Components and Plugins”.

In some respects, the server plugin API is similar to the older user-defined function (UDF) API that it supersedes, but the plugin API has several advantages over the older interface. For example, UDFs had no versioning information. Also, the newer plugin interface eliminates the security issues of the older UDF interface. The older interface for writing nonplugin UDFs permitted libraries to be loaded from any directory searched by the system's dynamic linker, and the symbols that identified the UDF library were relatively nonspecific.

The client plugin API has similar architectural characteristics, but client plugins have no direct access to the server the way server plugins do.

28.2.3 Plugin API Components

The server plugin implementation comprises several components.

SQL statements:

- `INSTALL PLUGIN` registers a plugin in the `mysql.plugin` table and loads the plugin code.
- `UNINSTALL PLUGIN` unregisters a plugin from the `mysql.plugin` table and unloads the plugin code.
- The `WITH PARSER` clause for full-text index creation associates a full-text parser plugin with a given `FULLTEXT` index.
- `SHOW PLUGINS` displays information about server plugins.

Command-line options and system variables:

- The `--plugin-load` option enables plugins to be loaded at server startup time.
- The `plugin_dir` system variable indicates the location of the directory where all plugins must be installed. The value of this variable can be specified at server startup with a `--plugin_dir=`*`dir_name`* option. `mysql_config --plugindir` displays the default plugin directory path name.

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

Plugin-related tables:

- The `INFORMATION_SCHEMA.PLUGINS` table contains plugin information.
- The `mysql.plugin` table lists each plugin that was installed with `INSTALL PLUGIN` and is required for plugin use. For new MySQL installations, this table is created during the installation process.

The client plugin implementation is simpler:

- For the `mysql_options()` C API function, the `MYSQL_DEFAULT_AUTH` and `MYSQL_PLUGIN_DIR` options enable client programs to load authentication plugins.
- There are C API functions that enable management of client plugins.

To examine how MySQL implements plugins, consult the following source files in a MySQL source distribution:

- In the `include/mysql` directory, `plugin.h` exposes the public plugin API. This file should be examined by anyone who wants to write a plugin library. `plugin_xxx.h` files provide additional information that pertains to specific types of plugins. `client_plugin.h` contains information specific to client plugins.
- In the `sql` directory, `sql_plugin.h` and `sql_plugin.cc` comprise the internal plugin implementation. `sql_acl.cc` is where the server uses authentication plugins. These files need not be consulted by plugin developers. They may be of interest for those who want to know more about how the server handles plugins.
- In the `sql-common` directory, `client_plugin.h` implements the C API client plugin functions, and `client.c` implements client authentication support. These files need not be consulted by plugin developers. They may be of interest for those who want to know more about how the server handles plugins.

28.2.4 Writing Plugins

To create a plugin library, you must provide the required descriptor information that indicates what plugins the library file contains, and write the interface functions for each plugin.

Every server plugin must have a general descriptor that provides information to the plugin API, and a type-specific descriptor that provides information about the plugin interface for a given type of plugin. The structure of the general descriptor is the same for all plugin types. The structure of the type-specific descriptor varies among plugin types and is determined by the requirements of what the plugin needs to do. The server plugin interface also enables plugins to expose status and system variables. These variables become visible through the `SHOW STATUS` and `SHOW VARIABLES` statements and the corresponding `INFORMATION_SCHEMA` tables.

For client-side plugins, the architecture is a bit different. Each plugin must have a descriptor, but there is no division into separate general and type-specific descriptors. Instead, the descriptor begins with a fixed set of members common to all client plugin types, and the common members are followed by any additional members required to implement the specific plugin type.

You can write plugins in C or C++ (or another language that can use C calling conventions). Plugins are loaded and unloaded dynamically, so your operating system must support dynamic loading and you must have compiled the calling application dynamically (not statically). For server plugins, this means that `mysqld` must be compiled dynamically.

A server plugin contains code that becomes part of the running server, so when you write the plugin, you are bound by any and all constraints that otherwise apply to writing server code. For example, you may have problems if you attempt to use functions from the `libstdc++` library. These constraints may change in future versions of the server, so it is possible that server upgrades will require revisions to plugins originally written for older servers. For information about these constraints, see Section 2.8.4, “MySQL Source-Configuration Options”, and Section 2.8.5, “Dealing with Problems Compiling MySQL”.

Client plugin writers should avoid dependencies on what symbols the calling application has because you cannot be sure what applications will use the plugin.

28.2.4.1 Overview of Plugin Writing

The following procedure provides an overview of the steps needed to create a plugin library. The next sections provide additional details on setting plugin data structures and writing specific types of plugins.

1. In the plugin source file, include the header files that the plugin library needs. The `plugin.h` file is required, and the library might require other files as well. For example:

```
#include <stdlib.h>
```

```
#include <ctype.h>
#include <mysql/plugin.h>
```

2. Set up the descriptor information for the plugin library file. For server plugins, write the library descriptor, which must contain the general plugin descriptor for each server plugin in the file. For more information, see Server Plugin Library and Plugin Descriptors. In addition, set up the type-specific descriptor for each server plugin in the library. Each plugin's general descriptor points to its type-specific descriptor.

   For client plugins, write the client descriptor. For more information, see Client Plugin Descriptors.

3. Write the plugin interface functions for each plugin. For example, each plugin's general plugin descriptor points to the initialization and deinitialization functions that the server should invoke when it loads and unloads the plugin. The plugin's type-specific description may also point to interface functions.

4. For server plugins, set up the status and system variables, if there are any.

5. Compile the plugin library as a shared library and install it in the plugin directory. For more information, see Section 28.2.4.3, “Compiling and Installing Plugin Libraries”.

6. For server plugins, register the plugin with the server. For more information, see Section 5.6.2, “Installing and Uninstalling Plugins”.

7. Test the plugin to verify that it works properly.

28.2.4.2 Plugin Data Structures

A plugin library file includes descriptor information to indicate what plugins it contains.

If the plugin library contains any server plugins, it must include the following descriptor information:

- A library descriptor indicates the general server plugin API version number used by the library and contains a general plugin descriptor for each server plugin in the library. To provide the framework for this descriptor, invoke two macros from the `plugin.h` header file:

  ```
  mysql_declare_plugin(name)
  ... one or more server plugin descriptors here ...
  mysql_declare_plugin_end;
  ```

  The macros expand to provide a declaration for the API version automatically. You must provide the plugin descriptors.

- Within the library descriptor, each general server plugin is described by a `st_mysql_plugin` structure. This plugin descriptor structure contains information that is common to every type of server plugin: A value that indicates the plugin type; the plugin name, author, description, and license type; pointers to the initialization and deinitialization functions that the server invokes when it loads and unloads the plugin, and pointers to any status or system variables the plugin implements.

- Each general server plugin descriptor within the library descriptor also contains a pointer to a type-specific plugin descriptor. The structure of the type-specific descriptors varies from one plugin type to another because each type of plugin can have its own API. A type-specific plugin descriptor contains a type-specific API version number and pointers to the functions that are needed to implement that plugin type. For example, a full-text parser plugin has initialization and deinitialization functions, and a main parsing function. The server invokes these functions when it uses the plugin to parse text.

The plugin library also contains the interface functions that are referenced by the general and type-specific descriptors for each plugin in the library.

If the plugin library contains a client plugin, it must include a descriptor for the plugin. The descriptor begins with a fixed set of members common to all client plugins, followed by any members specific to

the plugin type. To provide the descriptor framework, invoke two macros from the `client_plugin.h` header file:

```
mysql_declare_client_plugin(plugin_type)
   ... members common to all client plugins ...
   ... type-specific extra members ...
mysql_end_client_plugin;
```

The plugin library also contains any interface functions referenced by the client descriptor.

The `mysql_declare_plugin()` and `mysql_declare_client_plugin()` macros differ somewhat in how they can be invoked, which has implications for the contents of plugin libraries. The following guidelines summarize the rules:

- `mysql_declare_plugin()` and `mysql_declare_client_plugin()` can both be used in the same source file, which means that a plugin library can contain both server and client plugins. However, each of `mysql_declare_plugin()` and `mysql_declare_client_plugin()` can be used at most once.
- `mysql_declare_plugin()` permits multiple server plugin declarations, so a plugin library can contain multiple server plugins.
- `mysql_declare_client_plugin()` permits only a single client plugin declaration. To create multiple client plugins, separate plugin libraries must be used.

When a client program looks for a client plugin that is in a plugin library and not built into `libmysqlclient`, it looks for a file with a base name that is the same as the plugin name. For example, if a program needs to use a client authentication plugin named `auth_xxx` on a system that uses `.so` as the library suffix, it looks in the file named `auth_xxx.so`. (On OS X, the program looks first for `auth_xxx.dylib`, then for `auth_xxx.so`.) For this reason, if a plugin library contains a client plugin, the library must have the same base name as that plugin.

The same is not true for a library that contains server plugins. The `--plugin-load` option and the `INSTALL PLUGIN` statement provide the library file name explicitly, so there need be no explicit relationship between the library name and the name of any server plugins it contains.

Server Plugin Library and Plugin Descriptors

Every plugin library that contains server plugins must include a library descriptor that contains the general plugin descriptor for each server plugin in the file. This section discusses how to write the library and general descriptors for server plugins.

The library descriptor must define two symbols:

- `_mysql_plugin_interface_version_` specifies the version number of the general plugin framework. This is given by the `MYSQL_PLUGIN_INTERFACE_VERSION` symbol, which is defined in the `plugin.h` file.
- `_mysql_plugin_declarations_` defines an array of plugin declarations, terminated by a declaration with all members set to 0. Each declaration is an instance of the `st_mysql_plugin` structure (also defined in `plugin.h`). There must be one of these for each server plugin in the library.

If the server does not find those two symbols in a library, it does not accept it as a legal plugin library and rejects it with an error. This prevents use of a library for plugin purposes unless it was built specifically as a plugin library.

The conventional way to define the two required symbols is by using the `mysql_declare_plugin()` and `mysql_declare_plugin_end` macros from the `plugin.h` file:

```
mysql_declare_plugin(name)
 ... one or more server plugin descriptors here ...
mysql_declare_plugin_end;
```

Each server plugin must have a general descriptor that provides information to the server plugin API. The general descriptor has the same structure for all plugin types. The `st_mysql_plugin` structure in the `plugin.h` file defines this descriptor:

```
struct st_mysql_plugin
{
  int type;             /* the plugin type (a MYSQL_XXX_PLUGIN value)   */
  void *info;           /* pointer to type-specific plugin descriptor   */
  const char *name;     /* plugin name                                  */
  const char *author;   /* plugin author (for I_S.PLUGINS)              */
  const char *descr;    /* general descriptive text (for I_S.PLUGINS)   */
  int license;          /* the plugin license (PLUGIN_LICENSE_XXX)      */
  int (*init)(void *);  /* the function to invoke when plugin is loaded */
  int (*deinit)(void *);/* the function to invoke when plugin is unloaded */
  unsigned int version; /* plugin version (for I_S.PLUGINS)             */
  struct st_mysql_show_var *status_vars;
  struct st_mysql_sys_var **system_vars;
  void * __reserved1;   /* reserved for dependency checking             */
  unsigned long flags;  /* flags for plugin */
};
```

The `st_mysql_plugin` descriptor structure members are used as follows. `char *` members should be specified as null-terminated strings.

- `type`: The plugin type. This must be one of the plugin-type values from `plugin.h`:

  ```
  /*
    The allowable types of plugins
  */
  #define MYSQL_UDF_PLUGIN              0  /* User-defined function        */
  #define MYSQL_STORAGE_ENGINE_PLUGIN   1  /* Storage Engine               */
  #define MYSQL_FTPARSER_PLUGIN         2  /* Full-text parser plugin      */
  #define MYSQL_DAEMON_PLUGIN           3  /* The daemon/raw plugin type */
  #define MYSQL_INFORMATION_SCHEMA_PLUGIN  4  /* The I_S plugin type */
  #define MYSQL_AUDIT_PLUGIN            5  /* The Audit plugin type        */
  #define MYSQL_REPLICATION_PLUGIN      6  /* The replication plugin type */
  #define MYSQL_AUTHENTICATION_PLUGIN   7  /* The authentication plugin type */
  ...
  ```

  For example, for a full-text parser plugin, the `type` value is `MYSQL_FTPARSER_PLUGIN`.

- `info`: A pointer to the type-specific descriptor for the plugin. This descriptor's structure depends on the particular type of plugin, unlike that of the general plugin descriptor structure. For version-control purposes, the first member of the type-specific descriptor for every plugin type is expected to be the interface version for the type. This enables the server to check the type-specific version for every plugin no matter its type. Following the version number, the descriptor includes any other members needed, such as callback functions and other information needed by the server to invoke the plugin properly. Later sections on writing particular types of server plugins describe the structure of their type-specific descriptors.

- `name`: A string that gives the plugin name. This is the name that will be listed in the `mysql.plugin` table and by which you refer to the plugin in SQL statements such as `INSTALL PLUGIN` and `UNINSTALL PLUGIN`, or with the `--plugin-load` option. The name is also visible in the `INFORMATION_SCHEMA.PLUGINS` table or the output from `SHOW PLUGINS`.

  The plugin name should not begin with the name of any server option. If it does, the server will fail to initialize it. For example, the server has a `--socket` option, so you should not use a plugin name such as `socket`, `socket_plugin`, and so forth.

- `author`: A string naming the plugin author. This can be whatever you like.

- `desc`: A string that provides a general description of the plugin. This can be whatever you like.
- `license`: The plugin license type. The value can be one of `PLUGIN_LICENSE_PROPRIETARY`, `PLUGIN_LICENSE_GPL`, or `PLUGIN_LICENSE_BSD`.
- `init`: A once-only initialization function, or `NULL` if there is no such function. The server executes this function when it loads the plugin, which happens for `INSTALL PLUGIN` or, for plugins listed in the `mysql.plugin` table, at server startup. The function takes one argument that points to the internal structure used to identify the plugin. It returns zero for success and nonzero for failure.
- `deinit`: A once-only deinitialization function, or `NULL` if there is no such function. The server executes this function when it unloads the plugin, which happens for `UNINSTALL PLUGIN` or, for plugins listed in the `mysql.plugin` table, at server shutdown. The function takes one argument that points to the internal structure used to identify the plugin It returns zero for success and nonzero for failure.
- `version`: The plugin version number. When the plugin is installed, this value can be retrieved from the `INFORMATION_SCHEMA.PLUGINS` table. The value includes major and minor numbers. If you write the value as a hex constant, the format is `0x`*`MMNN`*, where *`MM`* and `NN` are the major and minor numbers, respectively. For example, `0x0302` represents version 3.2.
- `status_vars`: A pointer to a structure for status variables associated with the plugin, or `NULL` if there are no such variables. When the plugin is installed, these variables are displayed in the output of the `SHOW STATUS` statement.

  The `status_vars` member, if not `NULL`, points to an array of `st_mysql_show_var` structures that describe status variables. See Server Plugin Status and System Variables.
- `system_vars`: A pointer to a structure for system variables associated with the plugin, or `NULL` if there are no such variables. These options and system variables can be used to help initialize variables within the plugin.

  The `system_vars` member, if not `NULL`, points to an array of `st_mysql_sys_var` structures that describe system variables. See Server Plugin Status and System Variables.
- `__reserved1`: A placeholder for the future. It should be set to `NULL`.
- `flags`: Plugin flags. Individual bits correspond to different flags. The value should be set to the OR of the applicable flags. These flags are available:

  ```
  #define PLUGIN_OPT_NO_INSTALL   1UL   /* Not dynamically loadable */
  #define PLUGIN_OPT_NO_UNINSTALL 2UL   /* Not dynamically unloadable */
  ```

  `PLUGIN_OPT_NO_INSTALL` indicates that the plugin cannot be loaded at runtime with the `INSTALL PLUGIN` statement. This is appropriate for plugins that must be loaded at server startup with the `--plugin-load` option. `PLUGIN_OPT_NO_UNINSTALL` indicates that the plugin cannot be unloaded at runtime with the `UNINSTALL PLUGIN` statement.

The server invokes the `init` and `deinit` functions in the general plugin descriptor only when loading and unloading the plugin. They have nothing to do with use of the plugin such as happens when an SQL statement causes the plugin to be invoked.

For example, the descriptor information for a library that contains a single full-text parser plugin named `simple_parser` looks like this:

```
mysql_declare_plugin(ftexample)
{
  MYSQL_FTPARSER_PLUGIN,      /* type                            */
  &simple_parser_descriptor,  /* descriptor                      */
  "simple_parser",            /* name                            */
  "Oracle Corporation",       /* author                          */
```

```
  "Simple Full-Text Parser",  /* description                    */
  PLUGIN_LICENSE_GPL,         /* plugin license                 */
  simple_parser_plugin_init,  /* init function (when loaded)    */
  simple_parser_plugin_deinit,/* deinit function (when unloaded) */
  0x0001,                     /* version                        */
  simple_status,              /* status variables               */
  simple_system_variables,    /* system variables               */
  NULL,
  0
}
mysql_declare_plugin_end;
```

For a full-text parser plugin, the type must be `MYSQL_FTPARSER_PLUGIN`. This is the value that identifies the plugin as being legal for use in a `WITH PARSER` clause when creating a `FULLTEXT` index. (No other plugin type is legal for this clause.)

`plugin.h` defines the `mysql_declare_plugin()` and `mysql_declare_plugin_end` macros like this:

```
#ifndef MYSQL_DYNAMIC_PLUGIN
#define __MYSQL_DECLARE_PLUGIN(NAME, VERSION, PSIZE, DECLS) \
MYSQL_PLUGIN_EXPORT int VERSION= MYSQL_PLUGIN_INTERFACE_VERSION; \
MYSQL_PLUGIN_EXPORT int PSIZE= sizeof(struct st_mysql_plugin); \
MYSQL_PLUGIN_EXPORT struct st_mysql_plugin DECLS[]= {
#else
#define __MYSQL_DECLARE_PLUGIN(NAME, VERSION, PSIZE, DECLS) \
MYSQL_PLUGIN_EXPORT int _mysql_plugin_interface_version_= MYSQL_PLUGIN_INTERFACE_VERSION; \
MYSQL_PLUGIN_EXPORT int _mysql_sizeof_struct_st_plugin_= sizeof(struct st_mysql_plugin); \
MYSQL_PLUGIN_EXPORT struct st_mysql_plugin _mysql_plugin_declarations_[]= {
#endif

#define mysql_declare_plugin(NAME) \
__MYSQL_DECLARE_PLUGIN(NAME, \
                 builtin_ ## NAME ## _plugin_interface_version, \
                 builtin_ ## NAME ## _sizeof_struct_st_plugin, \
                 builtin_ ## NAME ## _plugin)

#define mysql_declare_plugin_end ,{0,0,0,0,0,0,0,0,0,0,0,0,0,0}}
```

Note

Those declarations define the `_mysql_plugin_interface_version_` symbol only if the `MYSQL_DYNAMIC_PLUGIN` symbol is defined. This means that `-DMYSQL_DYNAMIC_PLUGIN` must be provided as part of the compilation command to build the plugin as a shared library.

When the macros are used as just shown, they expand to the following code, which defines both of the required symbols (`_mysql_plugin_interface_version_` and `_mysql_plugin_declarations_`):

```
int _mysql_plugin_interface_version_= MYSQL_PLUGIN_INTERFACE_VERSION;
int _mysql_sizeof_struct_st_plugin_= sizeof(struct st_mysql_plugin);
struct st_mysql_plugin _mysql_plugin_declarations_[]= {
{
  MYSQL_FTPARSER_PLUGIN,      /* type                           */
  &simple_parser_descriptor,  /* descriptor                     */
  "simple_parser",            /* name                           */
  "Oracle Corporation",       /* author                         */
  "Simple Full-Text Parser",  /* description                    */
  PLUGIN_LICENSE_GPL,         /* plugin license                 */
  simple_parser_plugin_init,  /* init function (when loaded)    */
  simple_parser_plugin_deinit,/* deinit function (when unloaded) */
  0x0001,                     /* version                        */
  simple_status,              /* status variables               */
  simple_system_variables,    /* system variables               */
  NULL,
  0
```

```
}
  ,{0,0,0,0,0,0,0,0,0,0,0,0}}
};
```

The preceding example declares a single plugin in the general descriptor, but it is possible to declare multiple plugins. List the declarations one after the other between `mysql_declare_plugin()` and `mysql_declare_plugin_end`, separated by commas.

MySQL server plugins can be written in C or C++ (or another language that can use C calling conventions). If you write a C++ plugin, one C++ feature that you should not use is nonconstant variables to initialize global structures. Members of structures such as the `st_mysql_plugin` structure should be initialized only with constant variables. The `simple_parser` descriptor shown earlier is permissible in a C++ plugin because it satisfies that requirement:

```
mysql_declare_plugin(ftexample)
{
  MYSQL_FTPARSER_PLUGIN,      /* type                            */
  &simple_parser_descriptor,  /* descriptor                      */
  "simple_parser",            /* name                            */
  "Oracle Corporation",       /* author                          */
  "Simple Full-Text Parser",  /* description                     */
  PLUGIN_LICENSE_GPL,         /* plugin license                  */
  simple_parser_plugin_init,  /* init function (when loaded)     */
  simple_parser_plugin_deinit,/* deinit function (when unloaded) */
  0x0001,                     /* version                         */
  simple_status,              /* status variables                */
  simple_system_variables,    /* system variables                */
  NULL,
  0
}
mysql_declare_plugin_end;
```

Here is another valid way to write the general descriptor. It uses constant variables to indicate the plugin name, author, and description:

```
const char *simple_parser_name = "simple_parser";
const char *simple_parser_author = "Oracle Corporation";
const char *simple_parser_description = "Simple Full-Text Parser";

mysql_declare_plugin(ftexample)
{
  MYSQL_FTPARSER_PLUGIN,      /* type                            */
  &simple_parser_descriptor,  /* descriptor                      */
  simple_parser_name,         /* name                            */
  simple_parser_author,       /* author                          */
  simple_parser_description,  /* description                     */
  PLUGIN_LICENSE_GPL,         /* plugin license                  */
  simple_parser_plugin_init,  /* init function (when loaded)     */
  simple_parser_plugin_deinit,/* deinit function (when unloaded) */
  0x0001,                     /* version                         */
  simple_status,              /* status variables                */
  simple_system_variables,    /* system variables                */
  NULL,
  0
}
mysql_declare_plugin_end;
```

However, the following general descriptor is invalid. It uses structure members to indicate the plugin name, author, and description, but structures are not considered constant initializers in C++:

```
typedef struct
{
  const char *name;
  const char *author;
  const char *description;
} plugin_info;
```

```
plugin_info parser_info = {
  "simple_parser",
  "Oracle Corporation",
  "Simple Full-Text Parser"
};

mysql_declare_plugin(ftexample)
{
  MYSQL_FTPARSER_PLUGIN,      /* type                            */
  &simple_parser_descriptor,  /* descriptor                      */
  parser_info.name,           /* name                            */
  parser_info.author,         /* author                          */
  parser_info.description,    /* description                     */
  PLUGIN_LICENSE_GPL,         /* plugin license                  */
  simple_parser_plugin_init,  /* init function (when loaded)     */
  simple_parser_plugin_deinit,/* deinit function (when unloaded) */
  0x0001,                     /* version                         */
  simple_status,              /* status variables                */
  simple_system_variables,    /* system variables                */
  NULL,
  0
}
mysql_declare_plugin_end;
```

Server Plugin Status and System Variables

The server plugin interface enables plugins to expose status and system variables using the `status_vars` and `system_vars` members of the general plugin descriptor.

The `status_vars` member of the general plugin descriptor, if not 0, points to an array of `st_mysql_show_var` structures, each of which describes one status variable, followed by a structure with all members set to 0. The `st_mysql_show_var` structure has this definition:

```
struct st_mysql_show_var {
  const char *name;
  char *value;
  enum enum_mysql_show_type type;
};
```

When the plugin is installed, the plugin name and the `name` value are joined with an underscore to form the name displayed by `SHOW STATUS`.

The following table shows the permissible status variable `type` values and what the corresponding variable should be.

Table 28.1 Server Plugin Status Variable Types

Variable Type	Meaning
`SHOW_BOOL`	Pointer to a boolean variable
`SHOW_INT`	Pointer to an integer variable
`SHOW_LONG`	Pointer to a long integer variable
`SHOW_LONGLONG`	Pointer to a longlong integer variable
`SHOW_CHAR`	A string
`SHOW_CHAR_PTR`	Pointer to a string
`SHOW_ARRAY`	Pointer to another `st_mysql_show_var` array
`SHOW_FUNC`	Pointer to a function
`SHOW_DOUBLE`	Pointer to a double

For the `SHOW_FUNC` type, the function is called and fills in its `out` parameter, which then provides information about the variable to be displayed. The function has this signature:

```
#define SHOW_VAR_FUNC_BUFF_SIZE 1024

typedef int (*mysql_show_var_func) (void *thd,
                                    struct st_mysql_show_var *out,
                                    char *buf);
```

The `system_vars` member, if not 0, points to an array of `st_mysql_sys_var` structures, each of which describes one system variable (which can also be set from the command-line or configuration file), followed by a structure with all members set to 0. The `st_mysql_sys_var` structure is defined as follows:

```
struct st_mysql_sys_var {
 int flags;
 const char *name, *comment;
 int (*check)(THD*, struct st_mysql_sys_var *, void*, st_mysql_value*);
 void (*update)(THD*, struct st_mysql_sys_var *, void*, const void*);
};
```

Additional fields are append as required depending upon the flags.

For convenience, a number of macros are defined that make creating new system variables within a plugin much simpler.

Throughout the macros, the following fields are available:

- `name`: An unquoted identifier for the system variable.
- `varname`: The identifier for the static variable. Where not available, it is the same as the `name` field.
- `opt`: Additional use flags for the system variable. The following table shows the permissible flags.

  Table 28.2 Server Plugin System Variable Flags

Flag Value	Description
`PLUGIN_VAR_READONLY`	The system variable is read only
`PLUGIN_VAR_NOSYSVAR`	The system variable is not user visible at runtime
`PLUGIN_VAR_NOCMDOPT`	The system variable is not configurable from the command line
`PLUGIN_VAR_NOCMDARG`	No argument is required at the command line (typically used for boolean variables)
`PLUGIN_VAR_RQCMDARG`	An argument is required at the command line (this is the default)
`PLUGIN_VAR_OPCMDARG`	An argument is optional at the command line
`PLUGIN_VAR_MEMALLOC`	Used for string variables; indicates that memory is to be allocated for storage of the string

- `comment`: A descriptive comment to be displayed in the server help message. `NULL` if this variable is to be hidden.
- `check`: The check function, `NULL` for default.
- `update`: The update function, `NULL` for default.
- `default`: The variable default value.
- `minimum`: The variable minimum value.
- `maximum`: The variable maximum value.
- `blocksize`: The variable block size. When the value is set, it is rounded to the nearest multiple of `blocksize`.

A system variable may be accessed either by using the static variable directly or by using the `SYSVAR()`accessor macro. The `SYSVAR()` macro is provided for completeness. Usually it should be used only when the code cannot directly access the underlying variable.

For example:

```
static int my_foo;
static MYSQL_SYSVAR_INT(foo_var, my_foo,
                        PLUGIN_VAR_RQCMDARG, "foo comment",
                        NULL, NULL, 0, 0, INT_MAX, 0);
 ...
   SYSVAR(foo_var)= value;
   value= SYSVAR(foo_var);
   my_foo= value;
   value= my_foo;
```

Session variables may be accessed only through the `THDVAR()` accessor macro. For example:

```
static MYSQL_THDVAR_BOOL(some_flag,
                         PLUGIN_VAR_NOCMDARG, "flag comment",
                         NULL, NULL, FALSE);
 ...
   if (THDVAR(thd, some_flag))
   {
     do_something();
     THDVAR(thd, some_flag)= FALSE;
   }
```

All global and session system variables must be published to `mysqld` before use. This is done by constructing a `NULL`-terminated array of the variables and linking to it in the plugin public interface. For example:

```
static struct st_mysql_sys_var *my_plugin_vars[]= {
  MYSQL_SYSVAR(foo_var),
  MYSQL_SYSVAR(some_flag),
  NULL
};
mysql_declare_plugin(fooplug)
{
  MYSQL_..._PLUGIN,
  &plugin_data,
  "fooplug",
  "foo author",
  "This does foo!",
  PLUGIN_LICENSE_GPL,
  foo_init,
  foo_fini,
  0x0001,
  NULL,
  my_plugin_vars,
  NULL,
  0
}
mysql_declare_plugin_end;
```

The following convenience macros enable you to declare different types of system variables:

- Boolean system variables of type `bool`, which is a 1-byte boolean. (0 = `false`, 1 = `true`)

  ```
  MYSQL_THDVAR_BOOL(name, opt, comment, check, update, default)
  MYSQL_SYSVAR_BOOL(name, varname, opt, comment, check, update, default)
  ```

- String system variables of type `char*`, which is a pointer to a null-terminated string.

  ```
  ```

```
MYSQL_THDVAR_STR(name, opt, comment, check, update, default)
MYSQL_SYSVAR_STR(name, varname, opt, comment, check, update, default)
```

- Integer system variables, of which there are several varieties.
  - An `int` system variable, which is typically a 4-byte signed word.

```
MYSQL_THDVAR_INT(name, opt, comment, check, update, default, min, max, blk)
MYSQL_SYSVAR_INT(name, varname, opt, comment, check, update, default,
                 minimum, maximum, blocksize)
```

  - An `unsigned int` system variable, which is typically a 4-byte unsigned word.

```
MYSQL_THDVAR_UINT(name, opt, comment, check, update, default, min, max, blk)
MYSQL_SYSVAR_UINT(name, varname, opt, comment, check, update, default,
                  minimum, maximum, blocksize)
```

  - A `long` system variable, which is typically either a 4- or 8-byte signed word.

```
MYSQL_THDVAR_LONG(name, opt, comment, check, update, default, min, max, blk)
MYSQL_SYSVAR_LONG(name, varname, opt, comment, check, update, default,
                  minimum, maximum, blocksize)
```

  - An `unsigned long` system variable, which is typically either a 4- or 8-byte unsigned word.

```
MYSQL_THDVAR_ULONG(name, opt, comment, check, update, default, min, max, blk)
MYSQL_SYSVAR_ULONG(name, varname, opt, comment, check, update, default,
                   minimum, maximum, blocksize)
```

  - A `long long` system variable, which is typically an 8-byte signed word.

```
MYSQL_THDVAR_LONGLONG(name, opt, comment, check, update,
                      default, minimum, maximum, blocksize)
MYSQL_SYSVAR_LONGLONG(name, varname, opt, comment, check, update,
                      default, minimum, maximum, blocksize)
```

  - An `unsigned long long` system variable, which is typically an 8-byte unsigned word.

```
MYSQL_THDVAR_ULONGLONG(name, opt, comment, check, update,
                       default, minimum, maximum, blocksize)
MYSQL_SYSVAR_ULONGLONG(name, varname, opt, comment, check, update,
                       default, minimum, maximum, blocksize)
```

  - A `double` system variable, which is typically an 8-byte signed word.

```
MYSQL_THDVAR_DOUBLE(name, opt, comment, check, update,
                    default, minimum, maximum, blocksize)
MYSQL_SYSVAR_DOUBLE(name, varname, opt, comment, check, update,
                    default, minimum, maximum, blocksize)
```

  - An `unsigned long` system variable, which is typically either a 4- or 8-byte unsigned word. The range of possible values is an ordinal of the number of elements in the `typelib`, starting from 0.

```
MYSQL_THDVAR_ENUM(name, opt, comment, check, update, default, typelib)
MYSQL_SYSVAR_ENUM(name, varname, opt, comment, check, update,
                  default, typelib)
```

  - An `unsigned long long` system variable, which is typically an 8-byte unsigned word. Each bit represents an element in the `typelib`.

```
```

```
MYSQL_THDVAR_SET(name, opt, comment, check, update, default, typelib)
MYSQL_SYSVAR_SET(name, varname, opt, comment, check, update,
                 default, typelib)
```

Internally, all mutable and plugin system variables are stored in a `HASH` structure.

Display of the server command-line help text is handled by compiling a `DYNAMIC_ARRAY` of all variables relevant to command-line options, sorting them, and then iterating through them to display each option.

When a command-line option has been handled, it is then removed from the `argv` by the `handle_option()` function (`my_getopt.c`); in effect, it is consumed.

The server processes command-line options during the plugin installation process, immediately after the plugin has been successfully loaded but before the plugin initialization function has been called

Plugins loaded at runtime do not benefit from any configuration options and must have usable defaults. Once they are installed, they are loaded at `mysqld` initialization time and configuration options can be set at the command line or within `my.cnf`.

Plugins should consider the `thd` parameter to be read only.

Client Plugin Descriptors

Each client plugin must have a descriptor that provides information to the client plugin API. The descriptor structure begins with a fixed set of members common to all client plugins, followed by any members specific to the plugin type.

The `st_mysql_client_plugin` structure in the `client_plugin.h` file defines a “generic” descriptor that contains the common members:

```
struct st_mysql_client_plugin
{
  int type;
  unsigned int interface_version;
  const char *name;
  const char *author;
  const char *desc;
  unsigned int version[3];
  const char *license;
  void *mysql_api;
  int (*init)(char *, size_t, int, va_list);
  int (*deinit)();
  int (*options)(const char *option, const void *);
};
```

The common `st_mysql_client_plugin` descriptor structure members are used as follows. `char *` members should be specified as null-terminated strings.

- `type`: The plugin type. This must be one of the plugin-type values from `client_plugin.h`, such as `MYSQL_CLIENT_AUTHENTICATION_PLUGIN`.
- `interface_version`: The plugin interface version. For example, this is `MYSQL_CLIENT_AUTHENTICATION_PLUGIN_INTERFACE_VERSION` for an authentication plugin.
- `name`: A string that gives the plugin name. This is the name by which you refer to the plugin when you call `mysql_options()` with the `MYSQL_DEFAULT_AUTH` option or specify the `--default-auth` option to a MySQL client program.
- `author`: A string naming the plugin author. This can be whatever you like.
- `desc`: A string that provides a general description of the plugin. This can be whatever you like.

- `version`: The plugin version as an array of three integers indicating the major, minor, and teeny versions. For example, `{1,2,3}` indicates version 1.2.3.
- `license`: A string that specifies the license type.
- `mysql_api`: For internal use. Specify it as `NULL` in the plugin descriptor.
- `init`: A once-only initialization function, or `NULL` if there is no such function. The client library executes this function when it loads the plugin. The function returns zero for success and nonzero for failure.

  The `init` function uses its first two arguments to return an error message if an error occurs. The first argument is a pointer to a `char` buffer, and the second argument indicates the buffer length. Any message returned by the `init` function must be null-terminated, so the maximum message length is the buffer length minus one. The next arguments are passed to `mysql_load_plugin()`. The first indicates how many more arguments there are (0 if none), followed by any remaining arguments.
- `deinit`: A once-only deinitialization function, or `NULL` if there is no such function. The client library executes this function when it unloads the plugin. The function takes no arguments. It returns zero for success and nonzero for failure.
- `options`: A function for handling options passed to the plugin, or `NULL` if there is no such function. The function takes two arguments representing the option name and a pointer to its value. The function returns zero for success and nonzero for failure.

For a given client plugin type, the common descriptor members may be followed by additional members necessary to implement plugins of that type. For example, the `st_mysql_client_plugin_AUTHENTICATION` structure for authentication plugins has a function at the end that the client library calls to perform authentication.

To declare a plugin, use the `mysql_declare_client_plugin()` and `mysql_end_client_plugin` macros:

```
mysql_declare_client_plugin(plugin_type)
   ... members common to all client plugins ...
   ... type-specific extra members ...
mysql_end_client_plugin;
```

Do not specify the `type` or `interface_version` member explicitly. The `mysql_declare_client_plugin()` macro uses the *`plugin_type`* argument to generate their values automatically. For example, declare an authentication client plugin like this:

```
mysql_declare_client_plugin(AUTHENTICATION)
  "my_auth_plugin",
  "Author Name",
  "My Client Authentication Plugin",
  {1,0,0},
  "GPL",
  NULL,
  my_auth_init,
  my_auth_deinit,
  my_auth_options,
  my_auth_main
mysql_end_client_plugin;
```

This declaration uses the `AUTHENTICATION` argument to set the `type` and `interface_version` members to `MYSQL_CLIENT_AUTHENTICATION_PLUGIN` and `MYSQL_CLIENT_AUTHENTICATION_PLUGIN_INTERFACE_VERSION`.

Depending on the plugin type, the descriptor may have other members following the common members. For example, for an authentication plugin, there is a function (`my_auth_main()` in the descriptor just shown) that handles communication with the server. See Section 28.2.4.9, “Writing Authentication Plugins”.

Normally, a client program that supports the use of authentication plugins causes a plugin to be loaded by calling `mysql_options()` to set the `MYSQL_DEFAULT_AUTH` and `MYSQL_PLUGIN_DIR` options:

```
char *plugin_dir = "path_to_plugin_dir";
char *default_auth = "plugin_name";

/* ... process command-line options ... */

mysql_options(&mysql, MYSQL_PLUGIN_DIR, plugin_dir);
mysql_options(&mysql, MYSQL_DEFAULT_AUTH, default_auth);
```

Typically, the program will also accept `--plugin-dir` and `--default-auth` options that enable users to override the default values.

Should a client program require lower-level plugin management, the client library contains functions that take an `st_mysql_client_plugin` argument. See Section 27.7.13, “C API Client Plugin Functions”.

28.2.4.3 Compiling and Installing Plugin Libraries

After your plugin is written, you must compile it and install it. The procedure for compiling shared objects varies from system to system. If you build your library using `CMake`, it should be able to generate the correct compilation commands for your system. If the library is named `somepluglib`, you should end up with a shared library file that has a name something like `somepluglib.so`. (The `.so` file name suffix might differ on your system.)

To use `CMake`, you'll need to set up the configuration files to enable the plugin to be compiled and installed. Use the plugin examples under the `plugin` directory of a MySQL source distribution as a guide.

Create `CMakeLists.txt`, which should look something like this:

```
MYSQL_ADD_PLUGIN(somepluglib somepluglib.c
  MODULE_ONLY MODULE_OUTPUT_NAME "somepluglib")
```

When `CMake` generates the `Makefile`, it should take care of passing to the compilation command the `-DMYSQL_DYNAMIC_PLUGIN` flag, and passing to the linker the `-lmysqlservices` flag, which is needed to link in any functions from services provided through the plugin services interface. See Section 28.3, “MySQL Services for Components and Plugins”.

Run `CMake`, then run `make`:

```
shell> cmake .
shell> make
```

If you need to specify configuration options to `CMake`, see Section 2.8.4, “MySQL Source-Configuration Options”, for a list. For example, you might want to specify `CMAKE_INSTALL_PREFIX` to indicate the MySQL base directory under which the plugin should be installed. You can see what value to use for this option with `SHOW VARIABLES`:

```
mysql> SHOW VARIABLES LIKE 'basedir';
+---------------+------------------+
| Variable_name | Value            |
+---------------+------------------+
| base          | /usr/local/mysql |
+---------------+------------------+
```

The location of the plugin directory where you should install the library is given by the `plugin_dir` system variable. For example:

```
```

```
mysql> SHOW VARIABLES LIKE 'plugin_dir';
+---------------+---------------------------------+
| Variable_name | Value                           |
+---------------+---------------------------------+
| plugin_dir    | /usr/local/mysql/lib/mysql/plugin |
+---------------+---------------------------------+
```

To install the plugin library, use `make`:

```
shell> make install
```

Verify that `make install` installed the plugin library in the proper directory. After installing it, make sure that the library permissions permit it to be executed by the server.

28.2.4.4 Writing Full-Text Parser Plugins

MySQL supports server-side full-text parser plugins with `MyISAM` and `InnoDB`. For introductory information about full-text parser plugins, see Full-Text Parser Plugins.

A full-text parser plugin can be used to replace or modify the built-in full-text parser. This section describes how to write a full-text parser plugin named `simple_parser`. This plugin performs parsing based on simpler rules than those used by the MySQL built-in full-text parser: Words are nonempty runs of whitespace characters.

The instructions use the source code in the `plugin/fulltext` directory of MySQL source distributions, so change location into that directory. The following procedure describes how the plugin library is created:

1. To write a full-text parser plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

   ```
   #include <mysql/plugin.h>
   ```

   `plugin.h` defines the `MYSQL_FTPARSER_PLUGIN` server plugin type and the data structures needed to declare the plugin.

2. Set up the library descriptor for the plugin library file.

   This descriptor contains the general plugin descriptor for the server plugin. For a full-text parser plugin, the type must be `MYSQL_FTPARSER_PLUGIN`. This is the value that identifies the plugin as being legal for use in a `WITH PARSER` clause when creating a `FULLTEXT` index. (No other plugin type is legal for this clause.)

   For example, the library descriptor for a library that contains a single full-text parser plugin named `simple_parser` looks like this:

   ```
   mysql_declare_plugin(ftexample)
   {
     MYSQL_FTPARSER_PLUGIN,      /* type                            */
     &simple_parser_descriptor,  /* descriptor                      */
     "simple_parser",            /* name                            */
     "Oracle Corporation",       /* author                          */
     "Simple Full-Text Parser",  /* description                     */
     PLUGIN_LICENSE_GPL,         /* plugin license                  */
     simple_parser_plugin_init,  /* init function (when loaded)     */
     simple_parser_plugin_deinit,/* deinit function (when unloaded) */
     0x0001,                     /* version                         */
     simple_status,              /* status variables                */
     simple_system_variables,    /* system variables                */
     NULL,
     0
   ```

```
}
mysql_declare_plugin_end;
```

The `name` member (`simple_parser`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `SHOW PLUGINS` or `INFORMATION_SCHEMA.PLUGINS`.

For more information, see Server Plugin Library and Plugin Descriptors.

3. Set up the type-specific plugin descriptor.

   Each general plugin descriptor in the library descriptor points to a type-specific descriptor. For a full-text parser plugin, the type-specific descriptor is an instance of the `st_mysql_ftparser` structure in the `plugin.h` file:

   ```
   struct st_mysql_ftparser
   {
     int interface_version;
     int (*parse)(MYSQL_FTPARSER_PARAM *param);
     int (*init)(MYSQL_FTPARSER_PARAM *param);
     int (*deinit)(MYSQL_FTPARSER_PARAM *param);
   };
   ```

   As shown by the structure definition, the descriptor has an interface version number and contains pointers to three functions.

   The interface version number is specified using a symbol, which is in the form: `MYSQL_xxx_INTERFACE_VERSION`. For full-text parser plugins, the symbol is `MYSQL_FTPARSER_INTERFACE_VERSION`. In the source code, you will find the actual interface version number for the full-text parser plugin defined in `include/mysql/plugin_ftparser.h`. The current interface version number is `0x0101`.

   The `init` and `deinit` members should point to a function or be set to 0 if the function is not needed. The `parse` member must point to the function that performs the parsing.

   In the `simple_parser` declaration, that descriptor is indicated by `&simple_parser_descriptor`. The descriptor specifies the version number for the full-text plugin interface (as given by `MYSQL_FTPARSER_INTERFACE_VERSION`), and the plugin's parsing, initialization, and deinitialization functions:

   ```
   static struct st_mysql_ftparser simple_parser_descriptor=
   {
     MYSQL_FTPARSER_INTERFACE_VERSION, /* interface version      */
     simple_parser_parse,              /* parsing function       */
     simple_parser_init,               /* parser init function   */
     simple_parser_deinit              /* parser deinit function */
   };
   ```

   A full-text parser plugin is used in two different contexts, indexing and searching. In both contexts, the server calls the initialization and deinitialization functions at the beginning and end of processing each SQL statement that causes the plugin to be invoked. However, during statement processing, the server calls the main parsing function in context-specific fashion:

   - For indexing, the server calls the parser for each column value to be indexed.
   - For searching, the server calls the parser to parse the search string. The parser might also be called for rows processed by the statement. In natural language mode, there is no need for the server to call the parser. For boolean mode phrase searches or natural language searches with query expansion, the parser is used to parse column values for information that is not in the index. Also, if a boolean mode search is done for a column that has no `FULLTEXT` index, the built-in parser will be called. (Plugins are associated with specific indexes. If there is no index, no plugin is used.)

The plugin declaration in the general plugin descriptor has `init` and `deinit` members that point initialization and deinitialization functions, and so does the type-specific plugin descriptor to which it points. However, these pairs of functions have different purposes and are invoked for different reasons:

- For the plugin declaration in the general plugin descriptor, the initialization and deinitialization functions are invoked when the plugin is loaded and unloaded.
- For the type-specific plugin descriptor, the initialization and deinitialization functions are invoked per SQL statement for which the plugin is used.

Each interface function named in the plugin descriptor should return zero for success or nonzero for failure, and each of them receives an argument that points to a `MYSQL_FTPARSER_PARAM` structure containing the parsing context. The structure has this definition:

```
typedef struct st_mysql_ftparser_param
{
  int (*mysql_parse)(struct st_mysql_ftparser_param *,
                     char *doc, int doc_len);
  int (*mysql_add_word)(struct st_mysql_ftparser_param *,
                        char *word, int word_len,
                        MYSQL_FTPARSER_BOOLEAN_INFO *boolean_info);
  void *ftparser_state;
  void *mysql_ftparam;
  struct charset_info_st *cs;
  char *doc;
  int length;
  int flags;
  enum enum_ftparser_mode mode;
} MYSQL_FTPARSER_PARAM;
```

The structure members are used as follows:

- `mysql_parse`: A pointer to a callback function that invokes the server's built-in parser. Use this callback when the plugin acts as a front end to the built-in parser. That is, when the plugin parsing function is called, it should process the input to extract the text and pass the text to the `mysql_parse` callback.

  The first parameter for this callback function should be the `param` value itself:

  ```
  param->mysql_parse(param, ...);
  ```

  A front end plugin can extract text and pass it all at once to the built-in parser, or it can extract and pass text to the built-in parser a piece at a time. However, in this case, the built-in parser treats the pieces of text as though there are implicit word breaks between them.

- `mysql_add_word`: A pointer to a callback function that adds a word to a full-text index or to the list of search terms. Use this callback when the parser plugin replaces the built-in parser. That is, when the plugin parsing function is called, it should parse the input into words and invoke the `mysql_add_word` callback for each word.

  The first parameter for this callback function should be the `param` value itself:

  ```
  param->mysql_add_word(param, ...);
  ```

- `ftparser_state`: This is a generic pointer. The plugin can set it to point to information to be used internally for its own purposes.
- `mysql_ftparam`: This is set by the server. It is passed as the first argument to the `mysql_parse` or `mysql_add_word` callback.

- `cs`: A pointer to information about the character set of the text, or 0 if no information is available.
- `doc`: A pointer to the text to be parsed.
- `length`: The length of the text to be parsed, in bytes.
- `flags`: Parser flags. This is zero if there are no special flags. The only nonzero flag is `MYSQL_FTFLAGS_NEED_COPY`, which means that `mysql_add_word()` must save a copy of the word (that is, it cannot use a pointer to the word because the word is in a buffer that will be overwritten.)

  This flag might be set or reset by MySQL before calling the parser plugin, by the parser plugin itself, or by the `mysql_parse()` function.
- `mode`: The parsing mode. This value will be one of the following constants:
  - `MYSQL_FTPARSER_SIMPLE_MODE`: Parse in fast and simple mode, which is used for indexing and for natural language queries. The parser should pass to the server only those words that should be indexed. If the parser uses length limits or a stopword list to determine which words to ignore, it should not pass such words to the server.
  - `MYSQL_FTPARSER_WITH_STOPWORDS`: Parse in stopword mode. This is used in boolean searches for phrase matching. The parser should pass all words to the server, even stopwords or words that are outside any normal length limits.
  - `MYSQL_FTPARSER_FULL_BOOLEAN_INFO`: Parse in boolean mode. This is used for parsing boolean query strings. The parser should recognize not only words but also boolean-mode operators and pass them to the server as tokens using the `mysql_add_word` callback. To tell the server what kind of token is being passed, the plugin needs to fill in a `MYSQL_FTPARSER_BOOLEAN_INFO` structure and pass a pointer to it.

Note

For `MyISAM`, the stopword list and `ft_min_word_len` and `ft_max_word_len` are checked inside the tokenizer. For `InnoDB`, the stopword list and equivalent word length variable settings (`innodb_ft_min_token_size` and `innodb_ft_max_token_size`) are checked outside of the tokenizer. As a result, `InnoDB` plugin parsers do not need to check the stopword list, `innodb_ft_min_token_size`, or `innodb_ft_max_token_size`. Instead, it is recommended that all words be returned to `InnoDB`. However, if you want to check stopwords within your plugin parser, use `MYSQL_FTPARSER_SIMPLE_MODE`, which is for full-text search index and natural language search. For `MYSQL_FTPARSER_WITH_STOPWORDS` and `MYSQL_FTPARSER_FULL_BOOLEAN_INFO` modes, it is recommended that all words be returned to `InnoDB` including stopwords, in case of phrase searches.

If the parser is called in boolean mode, the `param->mode` value will be `MYSQL_FTPARSER_FULL_BOOLEAN_INFO`. The `MYSQL_FTPARSER_BOOLEAN_INFO` structure that the parser uses for passing token information to the server looks like this:

```
typedef struct st_mysql_ftparser_boolean_info
{
  enum enum_ft_token_type type;
  int yesno;
  int weight_adjust;
  char wasign;
  char trunc;
  int position;
```

```
  /* These are parser state and must be removed. */
  char prev;
  char *quot;
} MYSQL_FTPARSER_BOOLEAN_INFO;
```

The parser should fill in the structure members as follows:

- `type`: The token type. The following table shows the permissible types.

Table 28.3 Full-Text Parser Token Types

Token Value	Meaning
`FT_TOKEN_EOF`	End of data
`FT_TOKEN_WORD`	A regular word
`FT_TOKEN_LEFT_PAREN`	The beginning of a group or subexpression
`FT_TOKEN_RIGHT_PAREN`	The end of a group or subexpression
`FT_TOKEN_STOPWORD`	A stopword

- `yesno`: Whether the word must be present for a match to occur. 0 means that the word is optional but increases the match relevance if it is present. Values larger than 0 mean that the word must be present. Values smaller than 0 mean that the word must not be present.
- `weight_adjust`: A weighting factor that determines how much a match for the word counts. It can be used to increase or decrease the word's importance in relevance calculations. A value of zero indicates no weight adjustment. Values greater than or less than zero mean higher or lower weight, respectively. The examples at Section 12.9.2, “Boolean Full-Text Searches”, that use the `<` and `>` operators illustrate how weighting works.
- `wasign`: The sign of the weighting factor. A negative value acts like the `~` boolean-search operator, which causes the word's contribution to the relevance to be negative.
- `trunc`: Whether matching should be done as if the boolean-mode `*` truncation operator had been given.
- `position`: Start position of the word in the document, in bytes. Used by `InnoDB` full-text search. For existing plugins that are called in boolean mode, support must be added for the position member.

Plugins should not use the `prev` and `quot` members of the `MYSQL_FTPARSER_BOOLEAN_INFO` structure.

Note

The plugin parser framework does not support:

- The `@distance` boolean operator.
- A leading plus sign (`+`) or minus sign (`-`) boolean operator followed by a space and then a word (`'+ apple'` or `'- apple'`). The leading plus or minus sign must be directly adjacent to the word, for example: `'+apple'` or `'-apple'`.

For information about boolean full-text search operators, see Section 12.9.2, “Boolean Full-Text Searches”.

4. Set up the plugin interface functions.

The general plugin descriptor in the library descriptor names the initialization and deinitialization functions that the server should invoke when it loads and unloads the plugin. For `simple_parser`, these functions do nothing but return zero to indicate that they succeeded:

```
static int simple_parser_plugin_init(void *arg __attribute__((unused)))
{
  return(0);
}

static int simple_parser_plugin_deinit(void *arg __attribute__((unused)))
{
  return(0);
}
```

Because those functions do not actually do anything, you could omit them and specify 0 for each of them in the plugin declaration.

The type-specific plugin descriptor for `simple_parser` names the initialization, deinitialization, and parsing functions that the server invokes when the plugin is used. For `simple_parser`, the initialization and deinitialization functions do nothing:

```
static int simple_parser_init(MYSQL_FTPARSER_PARAM *param
                              __attribute__((unused)))
{
  return(0);
}

static int simple_parser_deinit(MYSQL_FTPARSER_PARAM *param
                                __attribute__((unused)))
{
  return(0);
}
```

Here too, because those functions do nothing, you could omit them and specify 0 for each of them in the plugin descriptor.

The main parsing function, `simple_parser_parse()`, acts as a replacement for the built-in full-text parser, so it needs to split text into words and pass each word to the server. The parsing function's first argument is a pointer to a structure that contains the parsing context. This structure has a `doc` member that points to the text to be parsed, and a `length` member that indicates how long the text is. The simple parsing done by the plugin considers nonempty runs of whitespace characters to be words, so it identifies words like this:

```
static int simple_parser_parse(MYSQL_FTPARSER_PARAM *param)
{
  char *end, *start, *docend= param->doc + param->length;

  for (end= start= param->doc;; end++)
  {
    if (end == docend)
    {
      if (end > start)
        add_word(param, start, end - start);
      break;
    }
    else if (isspace(*end))
    {
      if (end > start)
        add_word(param, start, end - start);
      start= end + 1;
    }
  }
  return(0);
}
```

As the parser finds each word, it invokes a function `add_word()` to pass the word to the server. `add_word()` is a helper function only; it is not part of the plugin interface. The parser passes the parsing context pointer to `add_word()`, as well as a pointer to the word and a length value:

```
static void add_word(MYSQL_FTPARSER_PARAM *param, char *word, size_t len)
{
  MYSQL_FTPARSER_BOOLEAN_INFO bool_info=
    { FT_TOKEN_WORD, 0, 0, 0, 0, 0, ' ', 0 };

  param->mysql_add_word(param, word, len, &bool_info);
}
```

For boolean-mode parsing, `add_word()` fills in the members of the `bool_info` structure as described earlier in the discussion of the `st_mysql_ftparser_boolean_info` structure.

5. Set up the status variables. For the `simple_parser` plugin, the following status variable array sets up one status variable with a value that is static text, and another with a value that is stored in a long integer variable:

```
long number_of_calls= 0;

struct st_mysql_show_var simple_status[]=
{
  {"static", (char *)"just a static text", SHOW_CHAR},
  {"called", (char *)&number_of_calls,     SHOW_LONG},
  {0,0,0}
};
```

When the plugin is installed, the plugin name and the `name` value are joined with an underscore to form the name displayed by `SHOW STATUS`. For the array just shown, the resulting status variable names are `simple_parser_static` and `simple_parser_called`. This convention means that you can easily display the variables for a plugin using its name:

```
mysql> SHOW STATUS LIKE 'simple_parser%';
+----------------------+--------------------+
| Variable_name        | Value              |
+----------------------+--------------------+
| simple_parser_static | just a static text |
| simple_parser_called | 0                  |
+----------------------+--------------------+
```

6. To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `simple_parser` plugin, it is compiled and installed when you build MySQL from source. It is also included in binary distributions. The build process produces a shared object library with a name of `mypluglib.so` (the `.so` suffix might differ depending on your platform).

7. To use the plugin, register it with the server. For example, to register the plugin at runtime, use this statement (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN simple_parser SONAME 'mypluglib.so';
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

8. To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement. See Section 5.6.3, “Obtaining Server Plugin Information”.

9. Test the plugin to verify that it works properly.

Create a table that contains a string column and associate the parser plugin with a `FULLTEXT` index on the column:

```
mysql> CREATE TABLE t (c VARCHAR(255),
    ->   FULLTEXT (c) WITH PARSER simple_parser
    -> ) ENGINE=MyISAM;
Query OK, 0 rows affected (0.01 sec)
```

Insert some text into the table and try some searches. These should verify that the parser plugin treats all nonwhitespace characters as word characters:

```
mysql> INSERT INTO t VALUES
    ->   ('utf8mb4_0900_as_cs is a case-sensitive collation'),
    ->   ('I\'d like a case of oranges'),
    ->   ('this is sensitive information'),
    ->   ('another row'),
    ->   ('yet another row');
Query OK, 5 rows affected (0.02 sec)
Records: 5  Duplicates: 0  Warnings: 0

mysql> SELECT c FROM t;
+-------------------------------------------------+
| c                                               |
+-------------------------------------------------+
| utf8mb4_0900_as_cs is a case-sensitive collation |
| I'd like a case of oranges                      |
| this is sensitive information                   |
| another row                                     |
| yet another row                                 |
+-------------------------------------------------+
5 rows in set (0.00 sec)

mysql> SELECT MATCH(c) AGAINST('case') FROM t;
+--------------------------+
| MATCH(c) AGAINST('case') |
+--------------------------+
|                        0 |
|          1.2968142032623 |
|                        0 |
|                        0 |
|                        0 |
+--------------------------+
5 rows in set (0.00 sec)

mysql> SELECT MATCH(c) AGAINST('sensitive') FROM t;
+-----------------------------+
| MATCH(c) AGAINST('sensitive') |
+-----------------------------+
|                           0 |
|                           0 |
|               1.3253291845322 |
|                           0 |
|                           0 |
+-----------------------------+
5 rows in set (0.01 sec)

mysql> SELECT MATCH(c) AGAINST('case-sensitive') FROM t;
+----------------------------------+
| MATCH(c) AGAINST('case-sensitive') |
+----------------------------------+
|                   1.3109166622162 |
|                                 0 |
|                                 0 |
|                                 0 |
|                                 0 |
+----------------------------------+
5 rows in set (0.01 sec)
```

```
mysql> SELECT MATCH(c) AGAINST('I\'d') FROM t;
+--------------------------+
| MATCH(c) AGAINST('I\'d') |
+--------------------------+
|                        0 |
|          1.2968142032623 |
|                        0 |
|                        0 |
|                        0 |
+--------------------------+
5 rows in set (0.01 sec)
```

Neither “case” nor “insensitive” match “case-insensitive” the way that they would for the built-in parser.

28.2.4.5 Writing Daemon Plugins

A daemon plugin is a simple type of plugin used for code that should be run by the server but that does not communicate with it. This section describes how to write a daemon server plugin, using the example plugin found in the `plugin/daemon_example` directory of MySQL source distributions. That directory contains the `daemon_example.cc` source file for a daemon plugin named `daemon_example` that writes a heartbeat string at regular intervals to a file named `mysql-heartbeat.log` in the data directory.

To write a daemon plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin.h>
```

`plugin.h` defines the `MYSQL_DAEMON_PLUGIN` server plugin type and the data structures needed to declare the plugin.

The `daemon_example.cc` file sets up the library descriptor as follows. The library descriptor includes a single general server plugin descriptor.

```
mysql_declare_plugin(daemon_example)
{
  MYSQL_DAEMON_PLUGIN,
  &daemon_example_plugin,
  "daemon_example",
  "Brian Aker",
  "Daemon example, creates a heartbeat beat file in mysql-heartbeat.log",
  PLUGIN_LICENSE_GPL,
  daemon_example_plugin_init, /* Plugin Init */
  daemon_example_plugin_deinit, /* Plugin Deinit */
  0x0100 /* 1.0 */,
  NULL,                       /* status variables                */
  NULL,                       /* system variables                */
  NULL,                       /* config options                  */
  0,                          /* flags                           */
}
mysql_declare_plugin_end;
```

The `name` member (`daemon_example`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `SHOW PLUGINS` or `INFORMATION_SCHEMA.PLUGINS`.

The second member of the plugin descriptor, `daemon_example_plugin`, points to the type-specific daemon plugin descriptor. This structure consists only of the type-specific API version number:

```
struct st_mysql_daemon daemon_example_plugin=
{ MYSQL_DAEMON_INTERFACE_VERSION  };
```

The type-specific structure has no interface functions. There is no communication between the server and the plugin, except that the server calls the initialization and deinitialization functions from the general plugin descriptor to start and stop the plugin:

- `daemon_example_plugin_init()` opens the heartbeat file and spawns a thread that wakes up periodically and writes the next message to the file.
- `daemon_example_plugin_deinit()` closes the file and performs other cleanup.

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `daemon_example` plugin, it is compiled and installed when you build MySQL from source. It is also included in binary distributions. The build process produces a shared object library with a name of `libdaemon_example.so` (the `.so` suffix might differ depending on your platform).

To use the plugin, register it with the server. For example, to register the plugin at runtime, use this statement (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN daemon_example SONAME 'libdaemon_example.so';
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement. See Section 5.6.3, “Obtaining Server Plugin Information”.

While the plugin is loaded, it writes a heartbeat string at regular intervals to a file named `mysql-heartbeat.log` in the data directory. This file grows without limit, so after you have satistifed yourself that the plugin operates correctly, unload it:

```
UNINSTALL PLUGIN daemon_example;
```

28.2.4.6 Writing INFORMATION_SCHEMA Plugins

This section describes how to write a server-side `INFORMATION_SCHEMA` table plugin. For example code that implements such plugins, see the `sql/sql_show.cc` file of a MySQL source distribution. You can also look at the example plugins found in the `InnoDB` source. See the `handler/i_s.cc` and `handler/ha_innodb.cc` files within the `InnoDB` source tree (in the `storage/innobase` directory).

To write an `INFORMATION_SCHEMA` table plugin, include the following header files in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <sql_class.h>
#include <table.h>
```

These header files are located in the `sql` directory of MySQL source distributions. They contain C++ structures, so the source file for an `INFORMATION_SCHEMA` plugin must be compiled as C++ (not C) code.

The source file for the example plugin developed here is named `simple_i_s_table.cc`. It creates a simple `INFORMATION_SCHEMA` table named `SIMPLE_I_S_TABLE` that has two columns named `NAME` and `VALUE`. The general descriptor for a plugin library that implements the table looks like this:

```
mysql_declare_plugin(simple_i_s_library)
{
  MYSQL_INFORMATION_SCHEMA_PLUGIN,
  &simple_table_info,               /* type-specific descriptor */
  "SIMPLE_I_S_TABLE",               /* table name */
  "Author Name",                    /* author */
```

```
  "Simple INFORMATION_SCHEMA table", /* description */
  PLUGIN_LICENSE_GPL,                /* license type */
  simple_table_init,                 /* init function */
  NULL,
  0x0100,                            /* version = 1.0 */
  NULL,                              /* no status variables */
  NULL,                              /* no system variables */
  NULL,                              /* no reserved information */
  0                                  /* no flags */
}
mysql_declare_plugin_end;
```

The `name` member (`SIMPLE_I_S_TABLE`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `SHOW PLUGINS` or `INFORMATION_SCHEMA.PLUGINS`.

The `simple_table_info` member of the general descriptor points to the type-specific descriptor, which consists only of the type-specific API version number:

```
static struct st_mysql_information_schema simple_table_info =
{ MYSQL_INFORMATION_SCHEMA_INTERFACE_VERSION };
```

The general descriptor points to the initialization and deinitialization functions:

- The initialization function provides information about the table structure and a function that populates the table.
- The deinitialization function performs any required cleanup. If no cleanup is needed, this descriptor member can be `NULL` (as in the example shown).

The initialization function should return 0 for success, 1 if an error occurs. The function receives a generic pointer, which it should interpret as a pointer to the table structure:

```
static int table_init(void *ptr)
{
  ST_SCHEMA_TABLE *schema_table= (ST_SCHEMA_TABLE*)ptr;

  schema_table->fields_info= simple_table_fields;
  schema_table->fill_table= simple_fill_table;
  return 0;
}
```

The function should set these two members of the table structure:

- `fields_info`: An array of `ST_FIELD_INFO` structures that contain information about each column.
- `fill_table`: A function that populates the table.

The array pointed to by `fields_info` should contain one element per column of the `INFORMATION_SCHEMA` plus a terminating element. The following `simple_table_fields` array for the example plugin indicates that `SIMPLE_I_S_TABLE` has two columns. `NAME` is string-valued with a length of 10 and `VALUE` is integer-valued with a display width of 20. The last structure marks the end of the array.

```
static ST_FIELD_INFO simple_table_fields[]=
{
  {"NAME", 10, MYSQL_TYPE_STRING, 0, 0 0, 0},
  {"VALUE", 6, MYSQL_TYPE_LONG, 0, MY_I_S_UNSIGNED, 0, 0},
  {0, 0, MYSQL_TYPE_NULL, 0, 0, 0, 0}
};
```

For more information about the column information structure, see the definition of `ST_FIELD_INFO` in the `table.h` header file. The permissible `MYSQL_TYPE_`*`xxx`* type values are those used in the C API; see Section 27.7.5, “C API Data Structures”.

The `fill_table` member should be set to a function that populates the table and returns 0 for success, 1 if an error occurs. For the example plugin, the `simple_fill_table()` function looks like this:

```
static int simple_fill_table(THD *thd, TABLE_LIST *tables, Item *cond)
{
  TABLE *table= tables->table;

  table->field[0]->store("Name 1", 6, system_charset_info);
  table->field[1]->store(1);
  if (schema_table_store_record(thd, table))
    return 1;
  table->field[0]->store("Name 2", 6, system_charset_info);
  table->field[1]->store(2);
  if (schema_table_store_record(thd, table))
    return 1;
  return 0;
}
```

For each row of the `INFORMATION_SCHEMA` table, this function initializes each column, then calls `schema_table_store_record()` to install the row. The `store()` method arguments depend on the type of value to be stored. For column 0 (`NAME`, a string), `store()` takes a pointer to a string, its length, and information about the character set of the string:

```
store(const char *to, uint length, CHARSET_INFO *cs);
```

For column 1 (`VALUE`, an integer), `store()` takes the value and a flag indicating whether it is unsigned:

```
store(longlong nr, bool unsigned_value);
```

For other examples of how to populate `INFORMATION_SCHEMA` tables, search for instances of `schema_table_store_record()` in `sql_show.cc`.

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable).

To test the plugin, install it:

```
mysql> INSTALL PLUGIN SIMPLE_I_S_TABLE SONAME 'simple_i_s_table.so';
```

Verify that the table is present:

```
mysql> SELECT TABLE_NAME FROM INFORMATION_SCHEMA.TABLES
    -> WHERE TABLE_NAME = 'SIMPLE_I_S_TABLE';
+------------------+
| TABLE_NAME       |
+------------------+
| SIMPLE_I_S_TABLE |
+------------------+
```

Try to select from it:

```
mysql> SELECT * FROM INFORMATION_SCHEMA.SIMPLE_I_S_TABLE;
+--------+-------+
| NAME   | VALUE |
+--------+-------+
| Name 1 |     1 |
| Name 2 |     2 |
+--------+-------+
```

Uninstall it:

```
mysql> UNINSTALL PLUGIN SIMPLE_I_S_TABLE;
```

28.2.4.7 Writing Semisynchronous Replication Plugins

This section describes how to write server-side semisynchronous replication plugins, using the example plugins found in the `plugin/semisync` directory of MySQL source distributions. That directory contains the source files for master and slave plugins named `rpl_semi_sync_master` and `rpl_semi_sync_slave`. The information here covers only how to set up the plugin framework. For details about how the plugins implement replication functions, see the source.

To write a semisynchronous replication plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin.h>
```

`plugin.h` defines the `MYSQL_REPLICATION_PLUGIN` server plugin type and the data structures needed to declare the plugin.

For the master side, `semisync_master_plugin.cc` contains this general descriptor for a plugin named `rpl_semi_sync_master`:

```
mysql_declare_plugin(semi_sync_master)
{
  MYSQL_REPLICATION_PLUGIN,
  &semi_sync_master_plugin,
  "rpl_semi_sync_master",
  "He Zhenxing",
  "Semi-synchronous replication master",
  PLUGIN_LICENSE_GPL,
  semi_sync_master_plugin_init, /* Plugin Init */
  semi_sync_master_plugin_deinit, /* Plugin Deinit */
  0x0100 /* 1.0 */,
  semi_sync_master_status_vars, /* status variables */
  semi_sync_master_system_vars, /* system variables */
  NULL,                         /* config options */
  0,                            /* flags */
}
mysql_declare_plugin_end;
```

For the slave side, `semisync_slave_plugin.cc` contains this general descriptor for a plugin named `rpl_semi_sync_slave`:

```
mysql_declare_plugin(semi_sync_slave)
{
  MYSQL_REPLICATION_PLUGIN,
  &semi_sync_slave_plugin,
  "rpl_semi_sync_slave",
  "He Zhenxing",
  "Semi-synchronous replication slave",
  PLUGIN_LICENSE_GPL,
  semi_sync_slave_plugin_init, /* Plugin Init */
  semi_sync_slave_plugin_deinit, /* Plugin Deinit */
  0x0100 /* 1.0 */,
  semi_sync_slave_status_vars,  /* status variables */
  semi_sync_slave_system_vars,  /* system variables */
  NULL,                         /* config options */
  0,                            /* flags */
}
mysql_declare_plugin_end;
```

For both the master and slave plugins, the general descriptor has pointers to the type-specific descriptor, the initialization and deinitialization functions, and to the status and system variables implemented by the plugin. For information about variable setup, see Server Plugin Status and

System Variables. The following remarks discuss the type-specific descriptor and the initialization and deinitialization functions for the master plugin but apply similarly to the slave plugin.

The `semi_sync_master_plugin` member of the master general descriptor points to the type-specific descriptor, which consists only of the type-specific API version number:

```
struct Mysql_replication semi_sync_master_plugin= {
  MYSQL_REPLICATION_INTERFACE_VERSION
};
```

The initialization and deinitialization function declarations look like this:

```
static int semi_sync_master_plugin_init(void *p);
static int semi_sync_master_plugin_deinit(void *p);
```

The initialization function uses the pointer to register transaction and binary logging “observers” with the server. After successful initialization, the server takes care of invoking the observers at the appropriate times. (For details on the observers, see the source files.) The deinitialization function cleans up by deregistering the observers. Each function returns 0 for success or 1 if an error occurs.

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `rpl_semi_sync_master` and `rpl_semi_sync_slave` plugins, they are compiled and installed when you build MySQL from source. They are also included in binary distributions. The build process produces shared object libraries with names of `semisync_master.so` and `semisync_slave.so` (the `.so` suffix might differ depending on your platform).

28.2.4.8 Writing Audit Plugins

This section describes how to write a server-side audit plugin, using the example plugin found in the `plugin/audit_null` directory of MySQL source distributions. The `audit_null.c` and `audit_null_variables.h` source files in that directory implement an audit plugin named `NULL_AUDIT`.

Note

Other examples of plugins that use the audit plugin API are the query rewrite plugin (see Section 5.6.4, “The Rewriter Query Rewrite Plugin”) and the Version Tokens plugin (see Section 5.6.5, “Version Tokens”).

Within the server, the pluggable audit interface is implemented in the `sql_audit.h` and `sql_audit.cc` files in the `sql` directory of MySQL source distributions. Additionally, several places in the server call the audit interface when an auditable event occurs, so that registered audit plugins can be notified about the event if necessary. To see where such calls occur, search the server source files for invocations of functions with names of the form `mysql_audit_xxx()`. Audit notification occurs for server operations such as these:

- Client connect and disconnect events
- Writing a message to the general query log (if the log is enabled)
- Writing a message to the error log
- Sending a query result to a client

To write an audit plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin_audit.h>
```

`plugin_audit.h` includes `plugin.h`, so you need not include the latter file explicitly. `plugin.h` defines the `MYSQL_AUDIT_PLUGIN` server plugin type and the data structures needed to declare the plugin. `plugin_audit.h` defines data structures specific to audit plugins.

Audit Plugin General Descriptor

An audit plugin, like any MySQL server plugin, has a general plugin descriptor (see Server Plugin Library and Plugin Descriptors) and a type-specific plugin descriptor. In `audit_null.c`, the general descriptor for `audit_null` looks like this:

```
mysql_declare_plugin(audit_null)
{
  MYSQL_AUDIT_PLUGIN,         /* type                            */
  &audit_null_descriptor,     /* descriptor                      */
  "NULL_AUDIT",               /* name                            */
  "Oracle Corp",              /* author                          */
  "Simple NULL Audit",        /* description                     */
  PLUGIN_LICENSE_GPL,
  audit_null_plugin_init,     /* init function (when loaded)     */
  audit_null_plugin_deinit,   /* deinit function (when unloaded) */
  0x0003,                     /* version                         */
  simple_status,              /* status variables                */
  system_variables,           /* system variables                */
  NULL,
  0,
}
mysql_declare_plugin_end;
```

The first member, `MYSQL_AUDIT_PLUGIN`, identifies this plugin as an audit plugin.

`audit_null_descriptor` points to the type-specific plugin descriptor, described later.

The `name` member (`NULL_AUDIT`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `INFORMATION_SCHEMA.PLUGINS` or `SHOW PLUGINS`.

The `audit_null_plugin_init` initialization function performs plugin initialization when the plugin is loaded. The `audit_null_plugin_deinit` function performs cleanup when the plugin is unloaded.

The general plugin descriptor also refers to `simple_status` and `system_variables`, structures that expose several status and system variables. When the plugin is enabled, these variables can be inspected using `SHOW` statements (`SHOW STATUS`, `SHOW VARIABLES`) or the appropriate Performance Schema tables.

The `simple_status` structure declares several status variables with names of the form `Audit_null_`*`xxx`*. `NULL_AUDIT` increments the `Audit_null_called` status variable for every notification that it receives. The other status variables are more specific and `NULL_AUDIT` increments them only for notifications of specific events.

`system_variables` is an array of system variable elements, each of which is defined using a `MYSQL_THDVAR_`*`xxx`* macro. These system variables have names of the form `null_audit_`*`xxx`*. These variables can be used to communicate with the plugin at runtime.

Audit Plugin Type-Specific Descriptor

The `audit_null_descriptor` value in the general plugin descriptor points to the type-specific plugin descriptor. For audit plugins, this descriptor has the following structure (defined in `plugin_audit.h`):

```
struct st_mysql_audit
{
  int interface_version;
  void (*release_thd)(MYSQL_THD);
  int (*event_notify)(MYSQL_THD, mysql_event_class_t, const void *);
  unsigned long class_mask[MYSQL_AUDIT_CLASS_MASK_SIZE];
```

```
};
```

The type-specific descriptor for audit plugins has these members:

- `interface_version`: By convention, type-specific plugin descriptors begin with the interface version for the given plugin type. The server checks `interface_version` when it loads the plugin to see whether the plugin is compatible with it. For audit plugins, the value of the `interface_version` member is `MYSQL_AUDIT_INTERFACE_VERSION` (defined in `plugin_audit.h`).
- `release_thd`: A function that the server calls to inform the plugin that it is being dissociated from its thread context. This should be `NULL` if there is no such function.
- `event_notify`: A function that the server calls to notify the plugin that an auditable event has occurred. This function should not be `NULL`; that would not make sense because no auditing would occur.
- `class_mask`: An array of `MYSQL_AUDIT_CLASS_MASK_SIZE` elements. Each element specifies a bitmask for a given event class to indicate the subclasses for which the plugin wants notification. (This is how the plugin “subscribes” to events of interest.) An element should be 0 to ignore events for the corresponding event class.

The server uses the `event_notify` and `release_thd` functions together. They are called within the context of a specific thread, and a thread might perform an activity that produces several event notifications. The first time the server calls `event_notify` for a thread, it creates a binding of the plugin to the thread. The plugin cannot be uninstalled while this binding exists. When no more events for the thread will occur, the server informs the plugin of this by calling the `release_thd` function, and then destroys the binding. For example, when a client issues a statement, the thread processing the statement might notify audit plugins about the result set produced by the statement and about the statement being logged. After these notifications occur, the server releases the plugin before putting the thread to sleep until the client issues another statement.

This design enables the plugin to allocate resources needed for a given thread in the first call to the `event_notify` function and release them in the `release_thd` function:

```
event_notify function:
  if memory is needed to service the thread
    allocate memory
  ... rest of notification processing ...

release_thd function:
  if memory was allocated
    release memory
  ... rest of release processing ...
```

That is more efficient than allocating and releasing memory repeatedly in the notification function.

For the `NULL_AUDIT` audit plugin, the type-specific plugin descriptor looks like this:

```
static struct st_mysql_audit audit_null_descriptor=
{
  MYSQL_AUDIT_INTERFACE_VERSION,                  /* interface version    */
  NULL,                                           /* release_thd function */
  audit_null_notify,                              /* notify function      */
  { (unsigned long) MYSQL_AUDIT_GENERAL_ALL,
    (unsigned long) MYSQL_AUDIT_CONNECTION_ALL,
    (unsigned long) MYSQL_AUDIT_PARSE_ALL,
    (unsigned long) MYSQL_AUDIT_AUTHORIZATION_ALL,
    (unsigned long) MYSQL_AUDIT_TABLE_ACCESS_ALL,
    (unsigned long) MYSQL_AUDIT_GLOBAL_VARIABLE_ALL,
    (unsigned long) MYSQL_AUDIT_SERVER_STARTUP_ALL,
    (unsigned long) MYSQL_AUDIT_SERVER_SHUTDOWN_ALL,
    (unsigned long) MYSQL_AUDIT_COMMAND_ALL,
    (unsigned long) MYSQL_AUDIT_QUERY_ALL,
```

```
    (unsigned long) MYSQL_AUDIT_STORED_PROGRAM_ALL }
};
```

The server calls `audit_null_notify()` to pass audit event information to the plugin. There is no `release_thd` function.

The `class_mask` member is an array that indicates which event classes the plugin subscribes to. As shown, the array contents subscribe to all subclasses of all event classes that are available. To ignore all notifications for a given event class, specify the corresponding `class_mask` element as 0.

The number of `class_mask` elements corresponds to the number of event classes, each of which is listed in the `mysql_event_class_t` enumeration defined in `plugin_audit.h`:

```
typedef enum
{
  MYSQL_AUDIT_GENERAL_CLASS          = 0,
  MYSQL_AUDIT_CONNECTION_CLASS       = 1,
  MYSQL_AUDIT_PARSE_CLASS            = 2,
  MYSQL_AUDIT_AUTHORIZATION_CLASS    = 3,
  MYSQL_AUDIT_TABLE_ACCESS_CLASS     = 4,
  MYSQL_AUDIT_GLOBAL_VARIABLE_CLASS  = 5,
  MYSQL_AUDIT_SERVER_STARTUP_CLASS   = 6,
  MYSQL_AUDIT_SERVER_SHUTDOWN_CLASS  = 7,
  MYSQL_AUDIT_COMMAND_CLASS          = 8,
  MYSQL_AUDIT_QUERY_CLASS            = 9,
  MYSQL_AUDIT_STORED_PROGRAM_CLASS   = 10,
  /* This item must be last in the list. */
  MYSQL_AUDIT_CLASS_MASK_SIZE
} mysql_event_class_t;
```

For any given event class, `plugin_audit.h` defines bitmask symbols for individual event subclasses, as well as an *`xxx`*`_ALL` symbol that is the union of the all subclass bitmasks. For example, for `MYSQL_AUDIT_CONNECTION_CLASS` (the class that covers connect and disconnect events), `plugin_audit.h` defines these symbols:

```
typedef enum
{
  /** occurs after authentication phase is completed. */
  MYSQL_AUDIT_CONNECTION_CONNECT          = 1 << 0,
  /** occurs after connection is terminated. */
  MYSQL_AUDIT_CONNECTION_DISCONNECT       = 1 << 1,
  /** occurs after COM_CHANGE_USER RPC is completed. */
  MYSQL_AUDIT_CONNECTION_CHANGE_USER      = 1 << 2,
  /** occurs before authentication. */
  MYSQL_AUDIT_CONNECTION_PRE_AUTHENTICATE = 1 << 3
} mysql_event_connection_subclass_t;

#define MYSQL_AUDIT_CONNECTION_ALL (MYSQL_AUDIT_CONNECTION_CONNECT | \
                                    MYSQL_AUDIT_CONNECTION_DISCONNECT | \
                                    MYSQL_AUDIT_CONNECTION_CHANGE_USER | \
                                    MYSQL_AUDIT_CONNECTION_PRE_AUTHENTICATE)
```

To subscribe to all subclasses of the connection event class (as the `NULL_AUDIT` plugin does), a plugin specifies `MYSQL_AUDIT_CONNECTION_ALL` in the corresponding `class_mask` element (`class_mask[1]` in this case). To subscribe to only some subclasses, the plugin sets the `class_mask` element to the union of the subclasses of interest. For example, to subscribe only to the connect and change-user subclasses, the plugin sets `class_mask[1]` to this value:

```
MYSQL_AUDIT_CONNECTION_CONNECT | MYSQL_AUDIT_CONNECTION_CHANGE_USER
```

Audit Plugin Notification Function

Most of the work for an audit plugin occurs in the notification function (the `event_notify` member of the type-specific plugin descriptor). The server calls this function for each auditable event. Audit plugin notification functions have this prototype:

```
int (*event_notify)(MYSQL_THD, mysql_event_class_t, const void *);
```

The second and third parameters of the `event_notify` function prototype represent the event class and a generic pointer to an event structure. (Events in different classes have different structures. The notification function can use the event class value to determine which event structure applies.) The function processes the event and returns a status indicating whether the server should continue processing the event or terminate it.

For `NULL_AUDIT`, the notification function is `audit_null_notify()`. This function increments a global event counter (which the plugin exposes as the value of the `Audit_null_called` status value), and then examines the event class to determine how to process the event structure:

```
static int audit_null_notify(MYSQL_THD thd __attribute__((unused)),
                             mysql_event_class_t event_class,
                             const void *event)
{
  ...

  number_of_calls++;

  if (event_class == MYSQL_AUDIT_GENERAL_CLASS)
  {
    const struct mysql_event_general *event_general=
                                    (const struct mysql_event_general *)event;
    ...
  }
  else if (event_class == MYSQL_AUDIT_CONNECTION_CLASS)
  {
    const struct mysql_event_connection *event_connection=
                                (const struct mysql_event_connection *) event;
    ...

  }
  else if (event_class == MYSQL_AUDIT_PARSE_CLASS)
  {
    const struct mysql_event_parse *event_parse =
                                      (const struct mysql_event_parse *)event;
    ...
  }
  ...
}
```

The notification function interprets the `event` argument according to the value of `event_class`. The `event` argument is a generic pointer to the event record, the structure of which differs per event class. (The `plugin_audit.h` file contains the structures that define the contents of each event class.) For each class, `audit_null_notify()` casts the event to the appropriate class-specific structure and then checks its subclass to determine which subclass counter to increment. For example, the code to handle events in the connection-event class looks like this:

```
else if (event_class == MYSQL_AUDIT_CONNECTION_CLASS)
{
  const struct mysql_event_connection *event_connection=
                               (const struct mysql_event_connection *) event;

  switch (event_connection->event_subclass)
  {
  case MYSQL_AUDIT_CONNECTION_CONNECT:
    number_of_calls_connection_connect++;
    break;
  case MYSQL_AUDIT_CONNECTION_DISCONNECT:
    number_of_calls_connection_disconnect++;
    break;
  case MYSQL_AUDIT_CONNECTION_CHANGE_USER:
    number_of_calls_connection_change_user++;
    break;
  case MYSQL_AUDIT_CONNECTION_PRE_AUTHENTICATE:
```

```
    number_of_calls_connection_pre_authenticate++;
      break;
  default:
    break;
  }
}
```

Note

The general event class (`MYSQL_AUDIT_GENERAL_CLASS`) is deprecated as of MySQL 5.7.9 and will be removed in a future MySQL release. To reduce plugin overhead, it is preferable to subscribe only to the more specific event classes of interest.

For some event classes, the `NULL_AUDIT` plugin performs other processing in addition to incrementing a counter. In any case, when the notification function finishes processing the event, it should return a status indicating whether the server should continue processing the event or terminate it.

Audit Plugin Error Handling

Audit plugin notification functions can report a status value for the current event two ways:

- Use the notification function return value. In this case, the function returns zero if the server should continue processing the event, or nonzero if the server should terminate the event.
- Call the `my_message()` function to set the error state before returning from the notification function. In this case, the notification function return value is ignored and the server terminates event processing with an error. The `my_message()` arguments indicate which error to report, and its message. For example:

  ```
  my_message(ER_AUDIT_API_ABORT, "This is my error message.", MYF(0));
  ```

  Some events cannot be aborted. A nonzero return value is not taken into consideration and the `my_message()` error call must follow an `is_error()` check. For example:

  ```
  if (!thd->get_stmt_da()->is_error())
  {
    my_message(ER_AUDIT_API_ABORT, "This is my error message.", MYF(0));
  }
  ```

Some events cannot be terminated:

- `MYSQL_AUDIT_CONNECTION_DISCONNECT`: The server cannot prevent a client from disconnecting.
- `MYSQL_AUDIT_COMMAND_END`: This event provides the status of a command that has finished executing, so there is no purpose to terminating it.

If an audit plugin returns nonzero status for a nonterminable event, the server ignores the status and continues processing the event. This is also true if an audit plugin uses the `my_message()` function to terminate a nonterminable event.

Audit Plugin Usage

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `NULL_AUDIT` plugin, it is compiled and installed when you build MySQL from source. It is also included in binary distributions. The build process produces a shared object library with a name of `adt_null.so` (the `.so` suffix might differ depending on your platform).

To register the plugin at runtime, use this statement (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN NULL_AUDIT SONAME 'adt_null.so';
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement. See Section 5.6.3, “Obtaining Server Plugin Information”.

While the audit plugin is installed, it exposes status variables that indicate the events for which the plugin has been called:

```
mysql> SHOW STATUS LIKE 'Audit_null%';
+---------------------------------------+--------+
| Variable_name                         | Value  |
+---------------------------------------+--------+
| Audit_null_authorization_column       | 0      |
| Audit_null_authorization_db           | 0      |
| Audit_null_authorization_procedure    | 0      |
| Audit_null_authorization_proxy        | 0      |
| Audit_null_authorization_table        | 0      |
| Audit_null_authorization_user         | 0      |
| Audit_null_called                     | 185547 |
| Audit_null_command_end                | 20999  |
| Audit_null_command_start              | 21001  |
| Audit_null_connection_change_user     | 0      |
| Audit_null_connection_connect         | 5823   |
| Audit_null_connection_disconnect      | 5818   |
| Audit_null_connection_pre_authenticate | 5823   |
| Audit_null_general_error              | 1      |
| Audit_null_general_log                | 26559  |
| Audit_null_general_result             | 19922  |
| Audit_null_general_status             | 21000  |
| Audit_null_global_variable_get        | 0      |
| Audit_null_global_variable_set        | 0      |
| Audit_null_parse_postparse            | 14648  |
| Audit_null_parse_preparse             | 14648  |
| Audit_null_query_nested_start         | 6      |
| Audit_null_query_nested_status_end    | 6      |
| Audit_null_query_start                | 14648  |
| Audit_null_query_status_end           | 14647  |
| Audit_null_server_shutdown            | 0      |
| Audit_null_server_startup             | 1      |
| Audit_null_table_access_delete        | 104    |
| Audit_null_table_access_insert        | 2839   |
| Audit_null_table_access_read          | 97842  |
| Audit_null_table_access_update        | 278    |
+---------------------------------------+--------+
```

`Audit_null_called` counts all events, and the other variables count instances of specific event subclasses. For example, the preceding `SHOW STATUS` statement causes the server to send a result to the client and to write a message to the general query log if that log is enabled. Thus, a client that issues the statement repeatedly causes `Audit_null_called`, `Audit_null_general_result`, and `Audit_null_general_log` to be incremented each time. Notifications occur whether or not that log is enabled.

The status variables values are aggregated across all sessions. There are no counters for individual sessions.

`NULL_AUDIT` exposes several system variables that enable communication with the plugin at runtime:

```
mysql> SHOW VARIABLES LIKE 'null_audit%';
+-----------------------------------+-------+
| Variable_name                     | Value |
+-----------------------------------+-------+
| null_audit_abort_message          |       |
| null_audit_abort_value            | 1     |
| null_audit_event_order_check      |       |
```

```
| null_audit_event_order_check_exact | 1     |
| null_audit_event_order_started     | 0     |
| null_audit_event_record            |       |
| null_audit_event_record_def        |       |
+------------------------------------+-------+
```

To check the order of audit API calls, set the `null_audit_event_order_check` variable to the expected event order. For example:

```
SET null_audit_event_order_check =
      'MYSQL_AUDIT_CONNECTION_PRE_AUTHENTICATE;;;'
      'MYSQL_AUDIT_GENERAL_LOG;;;'
      'MYSQL_AUDIT_CONNECTION_CONNECT;;';
```

The statement takes advantage of the SQL syntax that concatenates adjacent strings into a single string.

The format of the value is:

```
'event_name;event_data;command' [';event_name;event_data;command'] ...
```

After the event order is matched, the `null_audit_event_order_check` value is replaced with a value of `EVENT-ORDER-OK`.

Specifying a command value of `ABORT_RET` makes it possible to abort the audit API call on the specified event. The following example aborts `INSERT` statement execution when its `MYSQL_AUDIT_QUERY_STATUS_END` event occurs:

```
SET null_audit_event_order_check =
   'MYSQL_AUDIT_COMMAND_START;command_id="3";;'
   'MYSQL_AUDIT_GENERAL_LOG;;;'
   'MYSQL_AUDIT_QUERY_START;;;'
   'MYSQL_AUDIT_QUERY_STATUS_END;;ABORT_RET';
```

After the audit plugin matches the preceding sequence, it aborts event processing and sends an error message to the client:

```
ERROR 3164 (HY000): Aborted by Audit API ('MYSQL_AUDIT_QUERY_STATUS_END';1).
```

Returning a nonzero value from the audit API notification routine is the standard way to abort event execution. It is also possible to specify a custom error code by setting the `null_audit_abort_value` variable to the value that the notification routine should return:

```
SET null_audit_abort_value = 123;
```

Aborting a sequence results in a standard message with the custom error code. Suppose that you set audit log system variables like this:

```
SET null_audit_abort_value = 123;
SET null_audit_event_order_check =
    'MYSQL_AUDIT_COMMAND_START;command_id="3";;'
    'MYSQL_AUDIT_GENERAL_LOG;;;'
    'MYSQL_AUDIT_QUERY_START;;ABORT_RET';
```

Then execution of `SELECT 1` results in this error:

```
ERROR 3164 (HY000): Aborted by Audit API ('MYSQL_AUDIT_QUERY_START';123).
```

An event can be also aborted with a custom message, specified by setting the `null_audit_abort_message` variable: Suppose that you set audit log system variables like this:

```
SET null_audit_abort_message = 'Custom error text.';
SET null_audit_event_order_check =
    'MYSQL_AUDIT_COMMAND_START;command_id="3";;'
    'MYSQL_AUDIT_GENERAL_LOG;;;'
    'MYSQL_AUDIT_QUERY_START;;ABORT_RET';
```

Then aborting a sequence results in the following error:

```
ERROR 3164 (HY000): Custom error text.
```

For test-creation purposes, it is possible to record events that pass through the plugin. Recording starts by specifying start and end events in the `null_audit_event_record_def` variable:

```
SET null_audit_event_record_def =
    'MYSQL_AUDIT_COMMAND_START;MYSQL_AUDIT_COMMAND_END';
```

Statement execution results in storing the events that occur in the `null_audit_event_record` variable.

To disable the plugin after testing it, use this statement to unload it:

```
UNINSTALL PLUGIN NULL_AUDIT;
```

28.2.4.9 Writing Authentication Plugins

MySQL supports pluggable authentication, in which plugins are invoked to authenticate client connections. Authentication plugins enable the use of authentication methods other than the built-in method of passwords stored in the `mysql.user` table. For example, plugins can be written to access external authentication methods. Also, authentication plugins can support the proxy user capability, such that the connecting user is a proxy for another user and is treated, for purposes of access control, as having the privileges of a different user. For more information, see Section 6.3.10, “Pluggable Authentication”, and Section 6.3.11, “Proxy Users”.

An authentication plugin can be written for the server side or the client side. Server-side plugins use the same plugin API that is used for the other server plugin types such as full-text parser or audit plugins (although with a different type-specific descriptor). Client-side plugins use the client plugin API.

Several header files contain information relevant to authentication plugins:

- `plugin.h`: Defines the `MYSQL_AUTHENTICATION_PLUGIN` server plugin type.
- `client_plugin.h`: Defines the API for client plugins. This includes the client plugin descriptor and function prototypes for client plugin C API calls (see Section 27.7.13, “C API Client Plugin Functions”).
- `plugin_auth.h`: Defines the part of the server plugin API specific to authentication plugins. This includes the type-specific descriptor for server-side authentication plugins and the `MYSQL_SERVER_AUTH_INFO` structure.
- `plugin_auth_common.h`: Contains common elements of client and server authentication plugins. This includes return value definitions and the `MYSQL_PLUGIN_VIO` structure.

To write an authentication plugin, include the following header files in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

- For a source file that implements a server authentication plugin, include this file:

```
#include <mysql/plugin_auth.h>
```

- For a source file that implements a client authentication plugin, or both client and server plugins, include these files:

```
#include <mysql/plugin_auth.h>
#include <mysql/client_plugin.h>
#include <mysql.h>
```

`plugin_auth.h` includes `plugin.h` and `plugin_auth_common.h`, so you need not include the latter files explicitly.

This section describes how to write a pair of simple server and client authentication plugins that work together.

Warning

These plugins accept any non-empty password and the password is sent in clear text. This is insecure, so the plugins *should not be used in production environments.*

The server-side and client-side plugins developed here both are named `auth_simple`. As described in Section 28.2.4.2, “Plugin Data Structures”, the plugin library file must have the same base name as the client plugin, so the source file name is `auth_simple.c` and produces a library named `auth_simple.so` (assuming that your system uses `.so` as the suffix for library files).

In MySQL source distributions, authentication plugin source is located in the `plugin/auth` directory and can be examined as a guide to writing other authentication plugins. Also, to see how the built-in authentication plugins are implemented, see `sql/sql_acl.cc` for plugins that are built in to the MySQL server and `sql-common/client.c` for plugins that are built in to the `libmysqlclient` client library. (For the built-in client plugins, note that the `auth_plugin_t` structures used there differ from the structures used with the usual client plugin declaration macros. In particular, the first two members are provided explicitly, not by declaration macros.)

Writing the Server-Side Authentication Plugin

Declare the server-side plugin with the usual general descriptor format that is used for all server plugin types (see Server Plugin Library and Plugin Descriptors). For the `auth_simple` plugin, the descriptor looks like this:

```
mysql_declare_plugin(auth_simple)
{
  MYSQL_AUTHENTICATION_PLUGIN,
  &auth_simple_handler,                /* type-specific descriptor */
  "auth_simple",                       /* plugin name */
  "Author Name",                       /* author */
  "Any-password authentication plugin", /* description */
  PLUGIN_LICENSE_GPL,                  /* license type */
  NULL,                                /* no init function */
  NULL,                                /* no deinit function */
  0x0100,                              /* version = 1.0 */
  NULL,                                /* no status variables */
  NULL,                                /* no system variables */
  NULL,                                /* no reserved information */
  0                                    /* no flags */
}
mysql_declare_plugin_end;
```

The `name` member (`auth_simple`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `SHOW PLUGINS` or `INFORMATION_SCHEMA.PLUGINS`.

The `auth_simple_handler` member of the general descriptor points to the type-specific descriptor. For an authentication plugin, the type-specific descriptor is an instance of the `st_mysql_auth` structure (defined in `plugin_auth.h`):

```
struct st_mysql_auth
{
  int interface_version;
  const char *client_auth_plugin;
  int (*authenticate_user)(MYSQL_PLUGIN_VIO *vio, MYSQL_SERVER_AUTH_INFO *info);
  int (*generate_authentication_string)(char *outbuf,
      unsigned int *outbuflen, const char *inbuf, unsigned int inbuflen);
  int (*validate_authentication_string)(char* const inbuf, unsigned int buflen);
  int (*set_salt)(const char *password, unsigned int password_len,
                  unsigned char* salt, unsigned char *salt_len);
  const unsigned long authentication_flags;
};
```

The `st_mysql_auth` structure has these members:

- `interface_version`: The type-specific API version number, always `MYSQL_AUTHENTICATION_INTERFACE_VERSION`
- `client_auth_plugin`: The client plugin name
- `authenticate_user`: A pointer to the main plugin function that communicates with the client
- `generate_authentication_string`: A pointer to a plugin function that generates a password digest from an authentication string
- `validate_authentication_string`: A pointer to a plugin function that validates a password digest
- `set_salt`: A pointer to a plugin function that converts a scrambled password to binary form
- `authentication_flags`: A flags word

The `client_auth_plugin` member should indicate the name of the client plugin if a specific plugin is required. A value of `NULL` means “any plugin.” In the latter case, whatever plugin the client uses will do. This is useful if the server plugin does not care about the client plugin or what user name or password it sends. For example, this might be true if the server plugin authenticates only local clients and uses some property of the operating system rather than the information sent by the client plugin.

For `auth_simple`, the type-specific descriptor looks like this:

```
static struct st_mysql_auth auth_simple_handler =
{
  MYSQL_AUTHENTICATION_INTERFACE_VERSION,
  "auth_simple",             /* required client-side plugin name */
  auth_simple_server         /* server-side plugin main function */
  generate_auth_string_hash, /* generate digest from password string */
  validate_auth_string_hash, /* validate password digest */
  set_salt,                  /* generate password salt value */
  AUTH_FLAG_PRIVILEGED_USER_FOR_PASSWORD_CHANGE
};
```

The main function, `auth_simple_server()`, takes two arguments representing an I/O structure and a `MYSQL_SERVER_AUTH_INFO` structure. The structure definition, found in `plugin_auth.h`, looks like this:

```
typedef struct st_mysql_server_auth_info
{
  char *user_name;
  unsigned int user_name_length;
  const char *auth_string;
  unsigned long auth_string_length;
  char authenticated_as[MYSQL_USERNAME_LENGTH+1];
  char external_user[512];
  int  password_used;
  const char *host_or_ip;
```

```
  unsigned int host_or_ip_length;
} MYSQL_SERVER_AUTH_INFO;
```

The character set for string members is UTF-8. If there is a `_length` member associated with a string, it indicates the string length in bytes. Strings are also null-terminated.

When an authentication plugin is invoked by the server, it should interpret the `MYSQL_SERVER_AUTH_INFO` structure members as follows. Some of these are used to set the value of SQL functions or system variables within the client session, as indicated.

- `user_name`: The user name sent by the client. The value becomes the `USER()` function value.
- `user_name_length`: The length of `user_name` in bytes.
- `auth_string`: The value of the `authentication_string` column of the `mysql.user` table row for the matching account name (that is, the row that matches the client user name and host name and that the server uses to determine how to authenticate the client).

  Suppose that you create an account using the following statement:

  ```
  CREATE USER 'my_user'@'localhost'
    IDENTIFIED WITH my_plugin AS 'my_auth_string';
  ```

  When `my_user` connects from the local host, the server invokes `my_plugin` and passes `'my_auth_string'` to it as the `auth_string` value.
- `auth_string_length`: The length of `auth_string` in bytes.
- `authenticated_as`: The server sets this to the user name (the value of `user_name`). The plugin can alter it to indicate that the client should have the privileges of a different user. For example, if the plugin supports proxy users, the initial value is the name of the connecting (proxy) user, and the plugin can change this member to the proxied user name. The server then treats the proxy user as having the privileges of the proxied user (assuming that the other conditions for proxy user support are satisfied; see Implementing Proxy User Support in Authentication Plugins). The value is represented as a string at most `MYSQL_USER_NAME_LENGTH` bytes long, plus a terminating null. The value becomes the `CURRENT_USER()` function value.
- `external_user`: The server sets this to the empty string (null terminated). Its value becomes the `external_user` system variable value. If the plugin wants that system variable to have a different value, it should set this member accordingly; for example, to the connecting user name. The value is represented as a string at most 511 bytes long, plus a terminating null.
- `password_used`: This member applies when authentication fails. The plugin can set it or ignore it. The value is used to construct the failure error message of `Authentication fails. Password used: %s`. The value of `password_used` determines how `%s` is handled, as shown in the following table.

`password_used`	`%s` Handling
0	NO
1	YES
2	There will be no `%s`

- `host_or_ip`: The name of the client host if it can be resolved, or the IP address otherwise.
- `host_or_ip_length`: The length of `host_or_ip` in bytes.

The `auth_simple` main function, `auth_simple_server()`, reads the password (a null-terminated string) from the client and succeeds if the password is nonempty (first byte not null):

```
```

```
static int auth_simple_server (MYSQL_PLUGIN_VIO *vio,
                               MYSQL_SERVER_AUTH_INFO *info)
{
  unsigned char *pkt;
  int pkt_len;

  /* read the password as null-terminated string, fail on error */
  if ((pkt_len= vio->read_packet(vio, &pkt)) < 0)
    return CR_ERROR;

  /* fail on empty password */
  if (!pkt_len || *pkt == '\0')
  {
    info->password_used= PASSWORD_USED_NO;
    return CR_ERROR;
  }

  /* accept any nonempty password */
  info->password_used= PASSWORD_USED_YES;

  return CR_OK;
}
```

The main function should return one of the error codes shown in the following table.

Error Code	Meaning
`CR_OK`	Success
`CR_OK_HANDSHAKE_COMPLETE`	Do not send a status packet back to client
`CR_ERROR`	Error
`CR_AUTH_USER_CREDENTIALS`	Authentication failure
`CR_AUTH_HANDSHAKE`	Authentication handshake failure
`CR_AUTH_PLUGIN_ERROR`	Internal plugin error

For an example of how the handshake works, see the `plugin/auth/dialog.c` source file.

The server counts plugin errors in the Performance Schema `host_cache` table.

`auth_simple_server()` is so basic that it does not use the authentication information structure except to set the member that indicates whether a password was received.

A plugin that supports proxy users must return to the server the name of the proxied user (the MySQL user whose privileges the client user should get). To do this, the plugin must set the `info->authenticated_as` member to the proxied user name. For information about proxying, see Section 6.3.11, “Proxy Users”, and Implementing Proxy User Support in Authentication Plugins.

The `generate_authentication_string` member of the plugin descriptor takes the password and generates a password hash (digest) from it:

- The first two arguments are pointers to the output buffer and its maximum length in bytes. The function should write the password hash to the output buffer and reset the length to the actual hash length.
- The second two arguments indicate the password input buffer and its length in bytes.
- The function returns 0 for success, 1 if an error occurred.

For the `auth_simple` plugin, the `generate_auth_string_hash()` function implements the `generate_authentication_string` member. It just makes a copy of the password, unless it is too long to fit in the output buffer.

```
int generate_auth_string_hash(char *outbuf, unsigned int *buflen,
                              const char *inbuf, unsigned int inbuflen)
```

```
{
  /*
    fail if buffer specified by server cannot be copied to output buffer
  */
  if (*buflen < inbuflen)
    return 1;   /* error */
  strncpy(outbuf, inbuf, inbuflen);
  *buflen= strlen(inbuf);
  return 0;     /* success */
}
```

The `validate_authentication_string` member of the plugin descriptor validates a password hash:

- The arguments are a pointer to the password hash and its length in bytes.
- The function returns 0 for success, 1 if the password hash cannot be validated.

For the `auth_simple` plugin, the `validate_auth_string_hash()` function implements the `validate_authentication_string` member. It returns success unconditionally:

```
int validate_auth_string_hash(char* const inbuf  __attribute__((unused)),
                              unsigned int buflen  __attribute__((unused)))
{
  return 0;     /* success */
}
```

The `set_salt` member of the plugin descriptor is used only by the `mysql_native_password` plugin (see Section 6.5.1.1, “Native Pluggable Authentication”). For other authentication plugins, you can use this trivial implementation:

```
int set_salt(const char* password __attribute__((unused)),
             unsigned int password_len __attribute__((unused)),
             unsigned char* salt __attribute__((unused)),
             unsigned char* salt_len)
{
  *salt_len= 0;
  return 0;     /* success */
}
```

The `authentication_flags` member of the plugin descriptor contains flags that affect plugin operation. The permitted flags are:

- `AUTH_FLAG_PRIVILEGED_USER_FOR_PASSWORD_CHANGE`: Credential changes are a privileged operation. If this flag is set, the server requires that the user has the global `CREATE USER` privilege or the `UPDATE` privilege for the `mysql` database.
- `AUTH_FLAG_USES_INTERNAL_STORAGE`: Whether the plugin uses internal storage (in the `authentication_string` column of `mysql.user` rows). If this flag is not set, attempts to set the password using `SET PASSWORD` fail and the server produces a warning.

Writing the Client-Side Authentication Plugin

Declare the client-side plugin descriptor with the `mysql_declare_client_plugin()` and `mysql_end_client_plugin` macros (see Client Plugin Descriptors). For the `auth_simple` plugin, the descriptor looks like this:

```
mysql_declare_client_plugin(AUTHENTICATION)
  "auth_simple",                        /* plugin name */
  "Author Name",                        /* author */
  "Any-password authentication plugin", /* description */
  {1,0,0},                              /* version = 1.0.0 */
  "GPL",                                /* license type */
```

```
  NULL,                                 /* for internal use */
  NULL,                                 /* no init function */
  NULL,                                 /* no deinit function */
  NULL,                                 /* no option-handling function */
  auth_simple_client                    /* main function */
mysql_end_client_plugin;
```

The descriptor members from the plugin name through the option-handling function are common to all client plugin types. (For descriptions, see Client Plugin Descriptors.) Following the common members, the descriptor has an additional member specific to authentication plugins. This is the “main” function, which handles communication with the server. The function takes two arguments representing an I/O structure and a connection handler. For our simple any-password plugin, the main function does nothing but write to the server the password provided by the user:

```
static int auth_simple_client (MYSQL_PLUGIN_VIO *vio, MYSQL *mysql)
{
  int res;

  /* send password as null-terminated string in clear text */
  res= vio->write_packet(vio, (const unsigned char *) mysql->passwd,
                         strlen(mysql->passwd) + 1);

  return res ? CR_ERROR : CR_OK;
}
```

The main function should return one of the error codes shown in the following table.

Error Code	Meaning
CR_OK	Success
CR_OK_HANDSHAKE_COMPLETE	Success, client done
CR_ERROR	Error

CR_OK_HANDSHAKE_COMPLETE indicates that the client has done its part successfully and has read the last packet. A client plugin may return CR_OK_HANDSHAKE_COMPLETE if the number of round trips in the authentication protocol is not known in advance and the plugin must read another packet to determine whether authentication is finished.

Using the Authentication Plugins

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the plugin_dir system variable).

Register the server-side plugin with the server. For example, to load the plugin at server startup, use a --plugin-load=auth_simple.so option (adjust the .so suffix for your platform as necessary).

Create a user for whom the server will use the auth_simple plugin for authentication:

```
mysql> CREATE USER 'x'@'localhost'
    -> IDENTIFIED WITH auth_simple;
```

Use a client program to connect to the server as user x. The server-side auth_simple plugin communicates with the client program that it should use the client-side auth_simple plugin, and the latter sends the password to the server. The server plugin should reject connections that send an empty password and accept connections that send a nonempty password. Invoke the client program each way to verify this:

```
shell> mysql --user=x --skip-password
ERROR 1045 (28000): Access denied for user 'x'@'localhost' (using password: NO)

shell> mysql --user=x --password
```

```
Enter password: abc
mysql>
```

Because the server plugin accepts any nonempty password, it should be considered insecure. After testing the plugin to verify that it works, restart the server without the `--plugin-load` option so as not to indavertently leave the server running with an insecure authentication plugin loaded. Also, drop the user with `DROP USER 'x'@'localhost'`.

For additional information about loading and using authentication plugins, see Section 5.6.2, “Installing and Uninstalling Plugins”, and Section 6.3.10, “Pluggable Authentication”.

If you are writing a client program that supports the use of authentication plugins, normally such a program causes a plugin to be loaded by calling `mysql_options()` to set the `MYSQL_DEFAULT_AUTH` and `MYSQL_PLUGIN_DIR` options:

```
char *plugin_dir = "path_to_plugin_dir";
char *default_auth = "plugin_name";

/* ... process command-line options ... */

mysql_options(&mysql, MYSQL_PLUGIN_DIR, plugin_dir);
mysql_options(&mysql, MYSQL_DEFAULT_AUTH, default_auth);
```

Typically, the program will also accept `--plugin-dir` and `--default-auth` options that enable users to override the default values.

Should a client program require lower-level plugin management, the client library contains functions that take an `st_mysql_client_plugin` argument. See Section 27.7.13, “C API Client Plugin Functions”.

Implementing Proxy User Support in Authentication Plugins

One of the capabilities that pluggable authentication makes possible is proxy users (see Section 6.3.11, “Proxy Users”). For a server-side authentication plugin to participate in proxy user support, these conditions must be satisfied:

- When a connecting client should be treated as a proxy user, the plugin must return a different name in the `authenticated_as` member of the `MYSQL_SERVER_AUTH_INFO` structure, to indicate the proxied user name. It may also optionally set the `external_user` member, to set the value of the `external_user` system variable.
- Proxy user accounts must be set up to be authenticated by the plugin. Use the `CREATE USER` or `GRANT` statement to associate accounts with plugins.
- Proxy user accounts must have the `PROXY` privilege for the proxied accounts. Use the `GRANT` statement to grant this privilege.

In other words, the only aspect of proxy user support required of the plugin is that it set `authenticated_as` to the proxied user name. The rest is optional (setting `external_user`) or done by the DBA using SQL statements.

How does an authentication plugin determine which proxied user to return when the proxy user connects? That depends on the plugin. Typically, the plugin maps clients to proxied users based on the authentication string passed to it by the server. This string comes from the `AS` part of the `IDENTIFIED WITH` clause of the `CREATE USER` statement that specifies use of the plugin for authentication.

The plugin developer determines the syntax rules for the authentication string and implements the plugin according to those rules. Suppose that a plugin takes a comma-separated list of pairs that map external users to MySQL users. For example:

```
CREATE USER ''@'%.example.com'
```

```
  IDENTIFIED WITH my_plugin AS 'extuser1=mysqlusera, extuser2=mysqluserb'
CREATE USER ''@'%.example.org'
  IDENTIFIED WITH my_plugin AS 'extuser1=mysqluserc, extuser2=mysqluserd'
```

When the server invokes a plugin to authenticate a client, it passes the appropriate authentication string to the plugin. The plugin is responsible to:

1. Parse the string into its components to determine the mapping to use
2. Compare the client user name to the mapping
3. Return the proper MySQL user name

For example, if `extuser2` connects from an `example.com` host, the server passes `'extuser1=mysqlusera, extuser2=mysqluserb'` to the plugin, and the plugin should copy `mysqluserb` into `authenticated_as`, with a terminating null byte. If `extuser2` connects from an `example.org` host, the server passes `'extuser1=mysqluserc, extuser2=mysqluserd'`, and the plugin should copy `mysqluserd` instead.

If there is no match in the mapping, the action depends on the plugin. If a match is required, the plugin likely will return an error. Or the plugin might simply return the client name; in this case, it should not change `authenticated_as`, and the server will not treat the client as a proxy.

The following example demonstrates how to handle proxy users using a plugin named `auth_simple_proxy`. Like the `auth_simple` plugin described earlier, `auth_simple_proxy` accepts any nonempty password as valid (and thus should not be used in production environments). In addition, it examines the `auth_string` authentication string member and uses these very simple rules for interpreting it:

- If the string is empty, the plugin returns the user name as given and no proxying occurs. That is, the plugin leaves the value of `authenticated_as` unchanged.
- If the string is nonempty, the plugin treats it as the name of the proxied user and copies it to `authenticated_as` so that proxying occurs.

For testing, set up one account that is not proxied according to the preceding rules, and one that is. This means that one account has no `AS` clause, and one includes an `AS` clause that names the proxied user:

```
CREATE USER 'plugin_user1'@'localhost'
  IDENTIFIED WITH auth_simple_proxy;
CREATE USER 'plugin_user2'@'localhost'
  IDENTIFIED WITH auth_simple_proxy AS 'proxied_user';
```

In addition, create an account for the proxied user and grant `plugin_user2` the `PROXY` privilege for it:

```
CREATE USER 'proxied_user'@'localhost'
  IDENTIFIED BY 'proxied_user_pass';
GRANT PROXY
  ON 'proxied_user'@'localhost'
  TO 'plugin_user2'@'localhost';
```

Before the server invokes an authentication plugin, it sets `authenticated_as` to the client user name. To indicate that the user is a proxy, the plugin should set `authenticated_as` to the proxied user name. For `auth_simple_proxy`, this means that it must examine the `auth_string` value, and, if the value is nonempty, copy it to the `authenticated_as` member to return it as the name of the proxied user. In addition, when proxying occurs, the plugin sets the `external_user` member to the client user name; this becomes the value of the `external_user` system variable.

```
static int auth_simple_proxy_server (MYSQL_PLUGIN_VIO *vio,
                                     MYSQL_SERVER_AUTH_INFO *info)
{
```

```
  unsigned char *pkt;
  int pkt_len;

  /* read the password as null-terminated string, fail on error */
  if ((pkt_len= vio->read_packet(vio, &pkt)) < 0)
    return CR_ERROR;

  /* fail on empty password */
  if (!pkt_len || *pkt == '\0')
  {
    info->password_used= PASSWORD_USED_NO;
    return CR_ERROR;
  }

  /* accept any nonempty password */
  info->password_used= PASSWORD_USED_YES;

  /* if authentication string is nonempty, use as proxied user name */
  /* and use client name as external_user value */
  if (info->auth_string_length > 0)
  {
    strcpy (info->authenticated_as, info->auth_string);
    strcpy (info->external_user, info->user_name);
  }

  return CR_OK;
}
```

After a successful connection, the `USER()` function should indicate the connecting client user and host name, and `CURRENT_USER()` should indicate the account whose privileges apply during the session. The latter value should be the connecting user account if no proxying occurs or the proxied account if proxying does occur.

Compile and install the plugin, then test it. First, connect as `plugin_user1`:

```
shell> mysql --user=plugin_user1 --password
Enter password: x
```

In this case, there should be no proxying:

```
mysql> SELECT USER(), CURRENT_USER(), @@proxy_user, @@external_user\G
*************************** 1. row ***************************
        USER(): plugin_user1@localhost
 CURRENT_USER(): plugin_user1@localhost
   @@proxy_user: NULL
@@external_user: NULL
```

Then connect as `plugin_user2`:

```
shell> mysql --user=plugin_user2 --password
Enter password: x
```

In this case, `plugin_user2` should be proxied to `proxied_user`:

```
mysql> SELECT USER(), CURRENT_USER(), @@proxy_user, @@external_user\G
*************************** 1. row ***************************
        USER(): plugin_user2@localhost
 CURRENT_USER(): proxied_user@localhost
   @@proxy_user: 'plugin_user2'@'localhost'
@@external_user: 'plugin_user2'@'localhost'
```

28.2.4.10 Writing Password-Validation Plugins

This section describes how to write a server-side password-validation plugin. The instructions are based on the source code in the `plugin/password_validation` directory of MySQL source

distributions. The `validate_password.cc` source file in that directory implements the plugin named `validate_password`.

To write a password-validation plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin_validate_password.h>
```

`plugin_validate_password.h` includes `plugin.h`, so you need not include the latter file explicitly. `plugin.h` defines the `MYSQL_VALIDATE_PASSWORD_PLUGIN` server plugin type and the data structures needed to declare the plugin. `plugin_validate_password.h` defines data structures specific to password-validation plugins.

A password-validation plugin, like any MySQL server plugin, has a general plugin descriptor (see Server Plugin Library and Plugin Descriptors). In `validate_password.cc`, the general descriptor for `validate_password` looks like this:

```
mysql_declare_plugin(validate_password)
{
  MYSQL_VALIDATE_PASSWORD_PLUGIN,     /*   type                            */
  &validate_password_descriptor,      /*   descriptor                      */
  "validate_password",                /*   name                            */
  "Oracle Corporation",               /*   author                          */
  "check password strength",          /*   description                     */
  PLUGIN_LICENSE_GPL,
  validate_password_init,             /*   init function (when loaded)     */
  validate_password_deinit,           /*   deinit function (when unloaded) */
  0x0100,                             /*   version                         */
  NULL,
  validate_password_system_variables, /*   system variables                */
  NULL,
  0,
}
mysql_declare_plugin_end;
```

The `name` member (`validate_password`) indicates the name to use for references to the plugin in statements such as `INSTALL PLUGIN` or `UNINSTALL PLUGIN`. This is also the name displayed by `INFORMATION_SCHEMA.PLUGINS` or `SHOW PLUGINS`.

The general descriptor also refers to `validate_password_system_variables`, a structure that exposes several system variables to the `SHOW VARIABLES` statement:

```
static struct st_mysql_sys_var* validate_password_system_variables[]= {
  MYSQL_SYSVAR(length),
  MYSQL_SYSVAR(number_count),
  MYSQL_SYSVAR(mixed_case_count),
  MYSQL_SYSVAR(special_char_count),
  MYSQL_SYSVAR(policy),
  MYSQL_SYSVAR(dictionary_file),
  NULL
};
```

The `validate_password_init` initialization function reads the dictionary file if one was specified, and the `validate_password_deinit` function frees data structures associated with the file.

The `validate_password_descriptor` value in the general descriptor points to the type-specific descriptor. For password-validation plugins, this descriptor has the following structure:

```
struct st_mysql_validate_password
{
  int interface_version;
```

```
  /*
    This function returns TRUE for passwords which satisfy the password
    policy (as chosen by plugin variable) and FALSE for all other
    password
  */
  int (*validate_password)(mysql_string_handle password);
  /*
    This function returns the password strength (0-100) depending
    upon the policies
  */
  int (*get_password_strength)(mysql_string_handle password);
};
```

The type-specific descriptor has these members:

- `interface_version`: By convention, type-specific plugin descriptors begin with the interface version for the given plugin type. The server checks `interface_version` when it loads the plugin to see whether the plugin is compatible with it. For password-validation plugins, the value of the `interface_version` member is `MYSQL_VALIDATE_PASSWORD_INTERFACE_VERSION` (defined in `plugin_validate_password.h`).
- `validate_password`: A function that the server calls to test whether a password satisfies the current password policy. It returns 1 if the password is okay and 0 otherwise. The argument is the password, passed as a `mysql_string_handle` value. This data type is implemented by the `mysql_string` server service. For details, see the `string_service.h` and `string_service.cc` source files in the `sql` directory.
- `get_password_strength`: A function that the server calls to assess the strength of a password. It returns a value from 0 (weak) to 100 (strong). The argument is the password, passed as a `mysql_string_handle` value.

For the `validate_password` plugin, the type-specific descriptor looks like this:

```
static struct st_mysql_validate_password validate_password_descriptor=
{
  MYSQL_VALIDATE_PASSWORD_INTERFACE_VERSION,
  validate_password,                        /* validate function          */
  get_password_strength                     /* validate strength function */
};
```

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `validate_password` plugin, it is compiled and installed when you build MySQL from source. It is also included in binary distributions. The build process produces a shared object library with a name of `validate_password.so` (the `.so` suffix might differ depending on your platform).

To register the plugin at runtime, use this statement (adjust the `.so` suffix for your platform as necessary):

```
INSTALL PLUGIN validate_password SONAME 'validate_password.so';
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement. See Section 5.6.3, “Obtaining Server Plugin Information”.

While the `validate_password` plugin is installed, it exposes system variables that indicate the password-checking parameters:

```
mysql> SHOW VARIABLES LIKE 'validate_password%';
+--------------------------------------+--------+
```

```
| Variable_name                       | Value  |
+-------------------------------------+--------+
| validate_password_dictionary_file   |        |
| validate_password_length            | 8      |
| validate_password_mixed_case_count  | 1      |
| validate_password_number_count      | 1      |
| validate_password_policy            | MEDIUM |
| validate_password_special_char_count | 1      |
+-------------------------------------+--------+
```

For descriptions of these variables, see Section 6.5.3.2, “Password Validation Plugin Options and Variables”.

To disable the plugin after testing it, use this statement to unload it:

```
UNINSTALL PLUGIN validate_password;
```

28.2.4.11 Writing Protocol Trace Plugins

MySQL supports the use of protocol trace plugins: client-side plugins that implement tracing of communication between a client and the server that takes place using the client/server protocol.

Using the Test Protocol Trace Plugin

MySQL includes a test protocol trace plugin that serves to illustrate the information available from such plugins, and as a guide to writing other protocol trace plugins. To see how the test plugin works, use a MySQL source distribution; binary distributions are built with the test plugin disabled.

Enable the test protocol trace plugin by configuring MySQL with the `WITH_TEST_TRACE_PLUGIN` `CMake` option enabled. This causes the test trace plugin to be built and MySQL client programs to load it, but the plugin has no effect by default. Control the plugin using these environment variables:

- `MYSQL_TEST_TRACE_DEBUG`: Set this variable to a value other than 0 to cause the test plugin to produce diagnostic output on `stderr`.
- `MYSQL_TEST_TRACE_CRASH`: Set this variable to a value other than 0 to cause the test plugin to abort the client program if it detects an invalid trace event.

Caution

Diagnostic output from the test protocol trace plugin can disclose passwords and other sensitive information.

Given a MySQL installation built from source with the test plugin enabled, you can see a trace of the communication between the `mysql` client and the MySQL server as follows:

```
shell> export MYSQL_TEST_TRACE_DEBUG=1
shqll> mysql
test_trace: Test trace plugin initialized
test_trace: Starting tracing in stage CONNECTING
test_trace: stage: CONNECTING, event: CONNECTING
test_trace: stage: CONNECTING, event: CONNECTED
test_trace: stage: WAIT_FOR_INIT_PACKET, event: READ_PACKET
test_trace: stage: WAIT_FOR_INIT_PACKET, event: PACKET_RECEIVED
test_trace: packet received: 87 bytes
  0A 35 2E 37 2E 33 2D 6D  31 33 2D 64 65 62 75 67   .5.7.3-m13-debug
  2D 6C 6F 67 00 04 00 00  00 2B 7C 4F 55 3F 79 67   -log.....+|OU?yg
test_trace: 004: stage: WAIT_FOR_INIT_PACKET, event: INIT_PACKET_RECEIVED
test_trace: 004: stage: AUTHENTICATE, event: AUTH_PLUGIN
test_trace: 004: Using authentication plugin: mysql_native_password
test_trace: 004: stage: AUTHENTICATE, event: SEND_AUTH_RESPONSE
test_trace: 004: sending packet: 188 bytes
  85 A6 7F 00 00 00 00 01  21 00 00 00 00 00 00 00   .?......!.......
  00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00   ................
```

```
...
mysql> quit
test_trace: 008: stage: READY_FOR_COMMAND, event: SEND_COMMAND
test_trace: 008: QUIT
test_trace: 008: stage: READY_FOR_COMMAND, event: PACKET_SENT
test_trace: 008: packet sent: 0 bytes
test_trace: 008: stage: READY_FOR_COMMAND, event: DISCONNECTED
test_trace: 008: Connection  closed
test_trace: 008: Tracing connection has ended
Bye
test_trace: Test trace plugin de-initialized
```

To disable trace output, do this:

```
shell> MYSQL_TEST_TRACE_DEBUG=
```

Using Your Own Protocol Trace Plugins

Note

To use your own protocol trace plugins, you must configure MySQL with the `WITH_TEST_TRACE_PLUGIN` `CMake` option *disabled* because only one protocol trace plugin can be loaded at a time and an error occurs for attempts to load a second one. If you have already built MySQL with the test protocol trace plugin enabled to see how it works, you must rebuild MySQL without it before you can use your own plugins.

This section discusses how to write a basic protocol trace plugin named `simple_trace`. This plugin provides a framework showing how to set up the client plugin descriptor and create the trace-related callback functions. In `simple_trace`, these functions are rudimentary and do little other than illustrate the arguments required. To see in detail how a trace plugin can make use of trace event information, check the source file for the test protocol trace plugin (`test_trace_plugin.cc` in the `libmysql` directory of a MySQL source distribution). However, note that the `st_mysql_client_plugin_TRACE` structure used there differs from the structures used with the usual client plugin declaration macros. In particular, the first two members are defined explicitly, not implicitly by declaration macros.

Several header files contain information relevant to protocol trace plugins:

- `client_plugin.h`: Defines the API for client plugins. This includes the client plugin descriptor and function prototypes for client plugin C API calls (see Section 27.7.13, “C API Client Plugin Functions”).
- `plugin_trace.h`: Contains declarations for client-side plugins of type `MYSQL_CLIENT_TRACE_PLUGIN`. It also contains descriptions of the permitted protocol stages, transitions between stages, and the types of events permitted at each stage.

To write a protocol trace plugin, include the following header files in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin_trace.h>
#include <mysql.h>
```

`plugin_trace.h` includes `client_plugin.h`, so you need not include the latter file explicitly.

Declare the client-side plugin descriptor with the `mysql_declare_client_plugin()` and `mysql_end_client_plugin` macros (see Client Plugin Descriptors). For the `simple_trace` plugin, the descriptor looks like this:

```
```

```
mysql_declare_client_plugin(TRACE)
  "simple_trace",                 /* plugin name */
  "Author Name",                  /* author */
  "Simple protocol trace plugin", /* description */
  {1,0,0},                        /* version = 1.0.0 */
  "GPL",                          /* license type */
  NULL,                           /* for internal use */
  plugin_init,                    /* initialization function */
  plugin_deinit,                  /* deinitialization function */
  plugin_options,                 /* option-handling function */
  trace_start,                    /* start-trace function */
  trace_stop,                     /* stop-trace function */
  trace_event                     /* event-handling function */
mysql_end_client_plugin;
```

The descriptor members from the plugin name through the option-handling function are common to all client plugin types. The members following the common members implement trace event handling.

Function members for which the plugin needs no processing can be declared as `NULL` in the descriptor, in which case you need not write any corresponding function. For illustration purposes and to show the argument syntax, the following discussion implements all functions listed in the descriptor, even though some of them do nothing,

The initialization, deinitialization, and options functions common to all client plugins are declared as follows. For a description of the arguments and return values, see Client Plugin Descriptors.

```
static int
plugin_init(char *errbuf, size_t errbuf_len, int argc, va_list args)
{
  return 0;
}

static int
plugin_deinit()
{
  return 0;
}

static int
plugin_options(const char *option, const void *value)
{
  return 0;
}
```

The trace-specific members of the client plugin descriptor are callback functions. The following descriptions provide more detail on how they are used. Each has a first argument that is a pointer to the plugin instance in case your implementation needs to access it.

`trace_start()`: This function is called at the start of each traced connection (each connection that starts after the plugin is loaded). It is passed the connection handler and the protocol stage at which tracing starts. `trace_start()` allocates memory needed by the `trace_event()` function, if any, and returns a pointer to it. If no memory is needed, this function returns `NULL`.

```
static void*
trace_start(struct st_mysql_client_plugin_TRACE *self,
            MYSQL *conn,
            enum protocol_stage stage)
{
  struct st_trace_data *plugin_data= malloc(sizeof(struct st_trace_data));

  fprintf(stderr, "Initializing trace: stage %d\n", stage);
  if (plugin_data)
  {
    memset(plugin_data, 0, sizeof(struct st_trace_data));
    fprintf(stderr, "Trace initialized\n");
    return plugin_data;
```

```
  }
  fprintf(stderr, "Could not initialize trace\n");
  exit(1);
}
```

`trace_stop()`: This function is called when tracing of the connection ends. That usually happens when the connection is closed, but can happen earlier. For example, `trace_event()` can return a nonzero value at any time and that causes tracing of the connection to terminate. `trace_stop()` is then called even though the connection has not ended.

`trace_stop()` is passed the connection handler and a pointer to the memory allocated by `trace_start()` (`NULL` if none). If the pointer is non-`NULL`, `trace_stop()` should deallocate the memory. This function returns no value.

```
static void
trace_stop(struct st_mysql_client_plugin_TRACE *self,
           MYSQL *conn,
           void *plugin_data)
{
  fprintf(stderr, "Terminating trace\n");
  if (plugin_data)
    free(plugin_data);
}
```

`trace_event()`: This function is called for each event occurrence. It is passed a pointer to the memory allocated by `trace_start()` (`NULL` if none), the connection handler, the current protocol stage and event codes, and event data. This function returns 0 to continue tracing, nonzero if tracing should stop.

```
static int
trace_event(struct st_mysql_client_plugin_TRACE *self,
            void *plugin_data,
            MYSQL *conn,
            enum protocol_stage stage,
            enum trace_event event,
            struct st_trace_event_args args)
{
  fprintf(stderr, "Trace event received: stage %d, event %d\n", stage, event);
  if (event == TRACE_EVENT_DISCONNECTED)
    fprintf(stderr, "Connection closed\n");
  return 0;
}
```

The tracing framework shuts down tracing of the connection when the connection ends, so `trace_event()` should return nonzero only if you want to terminate tracing of the connection early. Suppose that you want to trace only connections for a certain MySQL account. After authentication, you can check the user name for the connection and stop tracing if it is not the user in whom you are interested.

For each call to `trace_event()`, the `st_trace_event_args` structure contains the event data. It has this definition:

```
struct st_trace_event_args
{
  const char           *plugin_name;
  int                   cmd;
  const unsigned char  *hdr;
  size_t                hdr_len;
  const unsigned char  *pkt;
  size_t                pkt_len;
};
```

For different event types, the `st_trace_event_args` structure contains the information described following. All lengths are in bytes. Unused members are set to `0`/`NULL`.

`AUTH_PLUGIN` event:

```
plugin_name  The name of the plugin
```

`SEND_COMMAND` event:

```
cmd          The command code
hdr          Pointer to the command packet header
hdr_len      Length of the header
pkt          Pointer to the command arguments
pkt_len      Length of the arguments
```

Other `SEND_xxx` and `xxx_RECEIVED` events:

```
pkt          Pointer to the data sent or received
pkt_len      Length of the data
```

`PACKET_SENT` event:

```
pkt_len      Number of bytes sent
```

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable).

After the plugin library file is compiled and installed in the plugin directory, you can test it easily by setting the `LIBMYSQL_PLUGINS` environment variable to the plugin name, which affects any client program that uses that variable. `mysql` is one such program:

```
shell> export LIBMYSQL_PLUGINS=simple_trace
shqll> mysql
Initializing trace: stage 0
Trace initialized
Trace event received: stage 0, event 1
Trace event received: stage 0, event 2
...
Welcome to the MySQL monitor.  Commands end with ; or \g.
Trace event received
Trace event received
...
mysql> SELECT 1;
Trace event received: stage 4, event 12
Trace event received: stage 4, event 16
...
Trace event received: stage 8, event 14
Trace event received: stage 8, event 15
+---+
| 1 |
+---+
| 1 |
+---+
1 row in set (0.00 sec)

mysql> quit
Trace event received: stage 4, event 12
Trace event received: stage 4, event 16
Trace event received: stage 4, event 3
Connection closed
Terminating trace
Bye
```

To stop the trace plugin from being loaded, do this:

```
shell> LIBMYSQL_PLUGINS=
```

It is also possible to write client programs that directly load the plugin. You can tell the client where the plugin directory is located by calling `mysql_options()` to set the `MYSQL_PLUGIN_DIR` option:

```
char *plugin_dir = "path_to_plugin_dir";

/* ... process command-line options ... */

mysql_options(&mysql, MYSQL_PLUGIN_DIR, plugin_dir);
```

Typically, the program will also accept a `--plugin-dir` option that enables users to override the default value.

Should a client program require lower-level plugin management, the client library contains functions that take an `st_mysql_client_plugin` argument. See Section 27.7.13, “C API Client Plugin Functions”.

28.2.4.12 Writing Keyring Plugins

MySQL Server supports a keyring service that enables internal server components and plugins to securely store sensitive information for later retrieval. This section describes how to write a server-side keyring plugin that can be used by service functions to perform key-management operations. For general keyring information, see Section 6.5.4, “The MySQL Keyring”.

The instructions here are based on the source code in the `plugin/keyring` directory of MySQL source distributions. The source files in that directory implement a plugin named `keyring_file` that uses a file local to the server host for data storage.

To write a keyring plugin, include the following header file in the plugin source file. Other MySQL or general header files might also be needed, depending on the plugin capabilities and requirements.

```
#include <mysql/plugin_keyring.h>
```

`plugin_keyring.h` includes `plugin.h`, so you need not include the latter file explicitly. `plugin.h` defines the `MYSQL_KEYRING_PLUGIN` server plugin type and the data structures needed to declare the plugin. `plugin_keyring.h` defines data structures specific to keyring plugins.

A keyring plugin, like any MySQL server plugin, has a general plugin descriptor (see Server Plugin Library and Plugin Descriptors). In `keyring.cc`, the general descriptor for `keyring_file` looks like this:

```
mysql_declare_plugin(keyring_file)
{
  MYSQL_KEYRING_PLUGIN,     /* type                                  */
  &keyring_descriptor,      /* descriptor                            */
  "keyring_file",           /* name                                  */
  "Oracle Corporation",     /* author                                */
  "store/fetch authentication data to/from a flat file", /* description */
  PLUGIN_LICENSE_GPL,
  keyring_init,             /* init function (when loaded)           */
  keyring_deinit,           /* deinit function (when unloaded)       */
  0x0100,                   /* version                               */
  NULL,                     /* status variables                      */
  keyring_system_variables, /* system variables                      */
  NULL,
  0,
}
mysql_declare_plugin_end;
```

The `name` member (`keyring_file`) indicates the plugin name. This is the name displayed by `INFORMATION_SCHEMA.PLUGINS` or `SHOW PLUGINS`.

The general descriptor also refers to `keyring_system_variables`, a structure that exposes a system variable to the `SHOW VARIABLES` statement:

```
static struct st_mysql_sys_var *keyring_system_variables[]= {
  MYSQL_SYSVAR(data),
  NULL
};
```

The `keyring_init` initialization function creates the data file if it does not exist, then reads it and initializes the keystore. The `keyring_deinit` function frees data structures associated with the file.

The `keyring_descriptor` value in the general descriptor points to the type-specific descriptor. For keyring plugins, this descriptor has the following structure:

```
struct st_mysql_keyring
{
  int interface_version;
  bool (*mysql_key_store)(const char *key_id, const char *key_type,
                          const char* user_id, const void *key, size_t key_len);
  bool (*mysql_key_fetch)(const char *key_id, char **key_type,
                          const char *user_id, void **key, size_t *key_len);
  bool (*mysql_key_remove)(const char *key_id, const char *user_id);
  bool (*mysql_key_generate)(const char *key_id, const char *key_type,
                             const char *user_id, size_t key_len);
};
```

The type-specific descriptor has these members:

- `interface_version`: By convention, type-specific plugin descriptors begin with the interface version for the given plugin type. The server checks `interface_version` when it loads the plugin to see whether the plugin is compatible with it. For keyring plugins, the value of the `interface_version` member is `MYSQL_KEYRING_INTERFACE_VERSION` (defined in `plugin_keyring.h`).
- `mysql_key_store`: A function that obfuscates and stores a key in the keyring.
- `mysql_key_fetch`: A function that deobfuscates and retrieves a key from the keyring.
- `mysql_key_remove`: A function that removes a key from the keyring.
- `mysql_key_generate`: A function that generates a new random key and stores it in the keyring.

For the `keyring_file` plugin, the type-specific descriptor looks like this:

```
static struct st_mysql_keyring keyring_descriptor=
{
  MYSQL_KEYRING_INTERFACE_VERSION,
  mysql_key_store,
  mysql_key_fetch,
  mysql_key_remove,
  mysql_key_generate
};
```

The `mysql_key_xxx` functions implemented by a keyring plugin are analogous to the `my_key_xxx` functions exposed by the keyring service API. For example, the `mysql_key_store` plugin function is analogous to the `my_key_store` keyring service function. For information about the arguments to keyring service functions and how they are used, see Section 28.3.2, “The Keyring Service”.

To compile and install a plugin library file, use the instructions in Section 28.2.4.3, “Compiling and Installing Plugin Libraries”. To make the library file available for use, install it in the plugin directory (the directory named by the `plugin_dir` system variable). For the `keyring_file` plugin, it is compiled and installed when you build MySQL from source. It is also included in binary distributions. The build

process produces a shared object library with a name of `keyring_file.so` (the `.so` suffix might differ depending on your platform).

Keyring plugins typically are loaded early during the server startup process so that they are available to built-in plugins and storage engines that might depend on them. For `keyring_file`, use these lines in the server `my.cnf` file (adjust the `.so` suffix for your platform as necessary):

```
[mysqld]
early-plugin-load=keyring_file.so
```

For additional information about plugin loading, see Section 5.6.2, “Installing and Uninstalling Plugins”.

To verify plugin installation, examine the `INFORMATION_SCHEMA.PLUGINS` table or use the `SHOW PLUGINS` statement (see Section 5.6.3, “Obtaining Server Plugin Information”). For example:

```
mysql> SELECT PLUGIN_NAME, PLUGIN_STATUS
       FROM INFORMATION_SCHEMA.PLUGINS
       WHERE PLUGIN_NAME LIKE 'keyring%';
+--------------+---------------+
| PLUGIN_NAME  | PLUGIN_STATUS |
+--------------+---------------+
| keyring_file | ACTIVE        |
+--------------+---------------+
```

While the `keyring_file` plugin is installed, it exposes a system variable that indicates the location of the data file it uses for secure information storage:

```
mysql> SHOW VARIABLES LIKE 'keyring_file%';
+-------------------+---------------------------------+
| Variable_name     | Value                           |
+-------------------+---------------------------------+
| keyring_file_data | /usr/local/mysql/keyring/keyring |
+-------------------+---------------------------------+
```

For a description of the `keyring_file_data` variable, see Section 5.1.5, “Server System Variables”.

To disable the plugin after testing it, restart the server without an `--early-plugin-load` option that names the plugin.

28.3 MySQL Services for Components and Plugins

MySQL server plugins have access to server “plugin services.” Similarly, server components have access to “component services.” The following discussion uses the term “plugins” but also applies to components.

The plugin services interface exposes server functionality that plugins can call. It complements the plugin API and has these characteristics:

- Services enable plugins to access code inside the server using ordinary function calls. Services are also available to user-defined functions (UDFs).
- Services are portable and work on multiple platforms.
- The interface includes a versioning mechanism so that service versions supported by the server can be checked at load time against plugin versions. Versioning protects against incompatibilities between the version of a service that the server provides and the version of the service expected or required by a plugin.
- For information about plugins for testing plugin services, see Plugins for Testing Plugin Services.

The plugin services interface differs from the plugin API as follows:

- The plugin API enables plugins to be used by the server. The calling initiative lies with the server to invoke plugins. This enables plugins to extend server functionality or register to receive notifications about server processing.
- The plugin services interface enables plugins to call code inside the server. The calling initiative lies with plugins to invoke service functions. This enables functionality already implemented in the server to be used by many plugins; they need not individually implement it themselves.

To determine what component services exist and what functions they provide, look in the `include/mysql/components` directory (and its `services` directory) of a MySQL source distribution.

To determine what plugin services exist and what functions they provide, look in the `include/mysql` directory of a MySQL source distribution. The relevant files are:

- `plugin.h` includes `services.h`, which is the "umbrella" header that includes all available service-specific header files.
- Service-specific headers have names of the form `service_xxx.h`.

The MySQL source code contains internal documentation written using Doxygen. This documentation is useful for understanding how MySQL works from a developer perspective. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/; in particular, see http://dev.mysql.com/doc/dev/mysql-server/latest/PAGE_EXTENDING.html. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, "Generating MySQL Doxygen Documentation Content".

Available component services include the following:

- `mysql_service_udf_registration`, `mysql_service_udf_registration_aggregate`: Services that enable components and plugins to register and unregister scalar and aggregate user-defined-functions (UDFs). Using these services, components and plugins can manage UDFs for themselves, without the need for `CREATE FUNCTION` and `DROP FUNCTION` statements. UDFs registered using these services are listed in the Performance Schema `user_defined_functions` table; see Section 25.11.16.4, "The user_defined_functions Table".

  These services were added in MySQL 8.0.2.

- `mysql_string`: A set of string service APIs.

  This service was converted from a plugin service to a component service in MySQL 8.0.2.

- `pfs_plugin_table`: A service for dynamic Performance Schema table manipulation.

  This service was added in MySQL 8.0.2.

Available plugin services include the following:

- `locking_service`: A service that implements locks with three attributes: Lock namespace, lock name, and lock mode. This locking interface is available at two levels: 1) As a C language interface, callable as a plugin service from server plugins or user-defined functions; 2) At the SQL level, as a set of user-defined functions that map onto calls to the service routines. For more information, see Section 28.3.1, "The Locking Service".
- `my_plugin_log_service`: A service that enables plugins to report errors and specify error messages. The server writes the messages to its error log.
- `my_snprintf`: A string-formatting service that produces consistent results across platforms.
- `my_thd_scheduler`: A service for plugins to select a thread scheduler.
- `mysql_keyring`: A service for keyring storage. For more information, see Section 28.3.2, "The Keyring Service".

- `mysql_password_policy`: A service for password validation and strength checking.
- `plugin_registry_service`: MySQL Server includes a component-based infrastructure for improving server extensibility; see Section 5.5, “MySQL Server Components”. However, MySQL plugins use an interface that predates the component interface. The `plugin_registry_service` enables plugins to access the component registry and its services.
- `security_context`: A service that enables plugins to examine or manipulate thread security contexts. This service provides setter and getter routines to access attributes of the server `Security_context` class, which includes attributes such as login user and host, authenticated user and host, and client IP address.
- `thd_alloc`: A memory-allocation service.
- `thd_wait`: A service for plugins to report when they are going to sleep or stall.

The remainder of this section describes how a plugin uses server functionality that is available as a service. See also the source for the “daemon” example plugin, which uses the `my_snprintf` service. Within a MySQL source distribution, that plugin is located in the `plugin/daemon_example` directory.

To use a service or services from within a plugin, the plugin source file must include the `plugin.h` header file to access service-related information:

```
#include <mysql/plugin.h>
```

This does not represent any additional setup cost. A plugin must include that file anyway because it contains definitions and structures that every plugin needs.

To access a service, a plugin calls service functions like any other function. For example, to format a string into a buffer for printing, call the `my_snprintf()` function provided by the service of the same name:

```
char buffer[BUFFER_SIZE];

my_snprintf(buffer, sizeof(buffer), format_string, argument_to_format, ...);
```

To report an error that the server will write to it error log, first choose an error level. `mysql/service_my_plugin_log.h` defines these levels:

```
enum plugin_log_level
{
  MY_ERROR_LEVEL,
  MY_WARNING_LEVEL,
  MY_INFORMATION_LEVEL
};
```

Then invoke `my_plugin_log_message()`:

```
int my_plugin_log_message(MYSQL_PLUGIN *plugin, enum plugin_log_level level,
                          const char *format, ...);
```

For example:

```
my_plugin_log_message(plugin_ptr, MY_ERROR_LEVEL, "Cannot initialize plugin");
```

Some services *for* plugins may be provided *by* plugins and thus are available only if the service-providing plugin is loaded. Any MySQL component that uses such a service should check whether the service is available.

When you build your plugin, use the `-lmysqlservices` flag at link time to link in the `libmysqlservices` library. For example, for `CMake`, put this in the top-level `CMakeLists.txt` file:

```
FIND_LIBRARY(MYSQLSERVICES_LIB mysqlservices
 PATHS "${MYSQL_SRCDIR}/libservices" NO_DEFAULT_PATH)
```

Put this in the `CMakeLists.txt` file in the directory containing the plugin source:

```
# the plugin needs the mysql services library for error logging
TARGET_LINK_LIBRARIES (your_plugin_library_name ${MYSQLSERVICES_LIB})
```

28.3.1 The Locking Service

Distributions provide a locking interface that is available at two levels:

- As a C language interface, callable as a plugin service from server plugins or user-defined functions
- At the SQL level, as a set of user-defined functions that map onto calls to the service routines

For general information about plugin services, see Section 28.3, “MySQL Services for Components and Plugins”. For general information about user-defined functions, see Section 28.4.2, “Adding a New User-Defined Function”.

The locking interface has these characteristics:

- Locks have three attributes: Lock namespace, lock name, and lock mode:
  - Locks are identified by the combination of namespace and lock name. The namespace enables different applications to use the same lock names without colliding by creating locks in separate namespaces. For example, if applications A and B use namespaces of `ns1` and `ns2`, respectively, each application can use lock names `lock1` and `lock2` without interfering with the other application.
  - A lock mode is either read or write. Read locks are shared: If a session has a read lock on a given lock identifier, other sessions can acquire a read lock on the same identifier. Write locks are exclusive: If a session has a write lock on a given lock identifier, other sessions cannot acquire a read or write lock on the same identifier.
- Namespace and lock names must be non-`NULL`, nonempty, and have a maximum length of 64 characters. A namespace or lock name specified as `NULL`, the empty string, or a string longer than 64 characters results in an `ER_LOCKING_SERVICE_WRONG_NAME` error.
- The locking interface treats namespace and lock names as binary strings, so comparisons are case sensitive.
- The locking interface provides functions to acquire locks and release locks. No special privilege is required to call these functions. Privilege checking is the responsibility of the calling application.
- Locks can be waited for if not immediately available. Lock acquisition calls take an integer timeout value that indicates how many seconds to wait to acquire locks before giving up. If the timeout is reached without successful lock acquisition, an `ER_LOCKING_SERVICE_TIMEOUT` error occurs. If the timeout is 0, there is no waiting and the call produces an error if locks cannot be acquired immediately.
- The locking interface detects deadlock between lock-acquisition calls in different sessions. In this case, the locking service chooses a caller and terminates its lock-acquisition request with an `ER_LOCKING_SERVICE_DEADLOCK` error. This error does not cause transactions to roll back. To choose a session in case of deadlock, the locking service prefers sessions that hold read locks over sessions that hold write locks.
- A session can acquire multiple locks with a single lock-acquisition call. For a given call, lock acquisition is atomic: The call succeeeds if all locks are acquired. If acquisition of any lock fails, the call acquires no locks and fails, typically with an `ER_LOCKING_SERVICE_TIMEOUT` or `ER_LOCKING_SERVICE_DEADLOCK` error.

- A session can acquire multiple locks for the same lock identifier (namespace and lock name combination). These lock instances can be read locks, write locks, or a mix of both.
- Locks acquired within a session are released explicitly by calling a release-locks function, or implicitly when the session terminates (either normally or abnormally). Locks are not released when transactions commit or roll back.
- Within a session, all locks for a given namespace when released are released together.

The interface provided by the locking service is distinct from that provided by `GET_LOCK()` and related SQL functions (see Section 12.20, “Miscellaneous Functions”). For example, `GET_LOCK()` does not implement namespaces and provides only exclusive locks, not distinct read and write locks.

28.3.1.1 The Locking Service C Interface

This section describes how to use the locking service C language interface. To use the UDF interface instead, see Section 28.3.1.2, “The Locking Service UDF Interface” For general characteristics of the locking service interface, see Section 28.3.1, “The Locking Service”. For general information about plugin services, see Section 28.3, “MySQL Services for Components and Plugins”.

Source files that use the locking service should include this header file:

```
#include <mysql/service_locking.h>
```

To acquire one or more locks, call this function:

```
int mysql_acquire_locking_service_locks(MYSQL_THD opaque_thd,
                                        const char* lock_namespace,
                                        const char**lock_names,
                                        size_t lock_num,
                                        enum enum_locking_service_lock_type lock_type,
                                        unsigned long lock_timeout);
```

The arguments have these meanings:

- `opaque_thd`: A thread handle. If specified as `NULL`, the handle for the current thread is used.
- `lock_namespace`: A null-terminated string that indicates the lock namespace.
- `lock_names`: An array of null-terminated strings that provides the names of the locks to acquire.
- `lock_num`: The number of names in the `lock_names` array.
- `lock_type`: The lock mode, either `LOCKING_SERVICE_READ` or `LOCKING_SERVICE_WRITE` to acquire read locks or write locks, respectively.
- `lock_timeout`: An integer number of seconds to wait to acquire the locks before giving up.

To release locks acquired for a given namespace, call this function:

```
int mysql_release_locking_service_locks(MYSQL_THD opaque_thd,
                                        const char* lock_namespace);
```

The arguments have these meanings:

- `opaque_thd`: A thread handle. If specified as `NULL`, the handle for the current thread is used.
- `lock_namespace`: A null-terminated string that indicates the lock namespace.

Locks acquired or waited for by the locking service can be monitored at the SQL level using the Performance Schema. For details, see Locking Service Monitoring.

28.3.1.2 The Locking Service UDF Interface

This section describes how to use the locking service user-defined function (UDF) interface. To use the C language interface instead, see Section 28.3.1.1, “The Locking Service C Interface” For general characteristics of the locking service interface, see Section 28.3.1, “The Locking Service”. For general information about user-defined functions, see Section 28.4.2, “Adding a New User-Defined Function”.

Installing or Uninstalling the UDF Locking Interface

The locking service routines described in Section 28.3.1.1, “The Locking Service C Interface” need not be installed because they are built into the server. The same is not true of the user-defined functions (UDFs) that map onto calls to the service routines: The UDFs must be installed before use. This section describes how to do that. For general information about UDF installation, see Section 28.4.2.5, “UDF Compiling and Installing”.

The locking service UDFs are implemented in a plugin library file located in the directory named by the `plugin_dir` system variable. The file base name is `locking_service`. The file name suffix differs per platform (for example, `.so` for Unix and Unix-like systems, `.dll` for Windows).

To install the locking service UDFs, use the `CREATE FUNCTION` statement (adjust the `.so` suffix for your platform as necessary):

```
CREATE FUNCTION service_get_read_locks RETURNS INT SONAME 'locking_service.so';
CREATE FUNCTION service_get_write_locks RETURNS INT SONAME 'locking_service.so';
CREATE FUNCTION service_release_locks RETURNS INT SONAME 'locking_service.so';
```

If the UDFs are used on a master replication server, install them on all slave servers as well to avoid replication problems.

Once installed, the UDFs remain installed until uninstalled. To remove them, use the `DROP FUNCTION` statement:

```
DROP FUNCTION service_get_read_locks;
DROP FUNCTION service_get_write_locks;
DROP FUNCTION service_release_locks;
```

Using the UDF Locking Interface

Before using the locking service UDFs, install them according to the instructions provided at Installing or Uninstalling the UDF Locking Interface.

To acquire one or more read locks, call this function:

```
mysql> SELECT service_get_read_locks('mynamespace', 'rlock1', 'rlock2', 10);
+-------------------------------------------------------------+
| service_get_read_locks('mynamespace', 'rlock1', 'rlock2', 10) |
+-------------------------------------------------------------+
|                                                           1 |
+-------------------------------------------------------------+
```

The first argument is the lock namespace. The final argument is an integer timeout indicating how many seconds to wait to acquire the locks before giving up. The arguments in between are the lock names.

For the example just shown, the function acquires locks with lock identifiers `(mynamespace, rlock1)` and `(mynamespace, rlock2)`.

To acquire write locks rather than read locks, call this function:

```
mysql> SELECT service_get_write_locks('mynamespace', 'wlock1', 'wlock2', 10);
+-------------------------------------------------------------+
```

```
| service_get_write_locks('mynamespace', 'wlock1', 'wlock2', 10) |
+----------------------------------------------------------------+
|                                                              1 |
+----------------------------------------------------------------+
```

In this case, the lock identifiers are `(mynamespace, wlock1)` and `(mynamespace, wlock2)`.

To release all locks for a namespace, use this function:

```
mysql> SELECT service_release_locks('mynamespace');
+--------------------------------------+
| service_release_locks('mynamespace') |
+--------------------------------------+
|                                    1 |
+--------------------------------------+
```

Each locking function returns nonzero for success. If the function fails, an error occurs. For example, the following error occurs because lock names cannot be empty:

```
mysql> SELECT service_get_read_locks('mynamespace', '', 10);
ERROR 3131 (42000): Incorrect locking service lock name ''.
```

A session can acquire multiple locks for the same lock identifier. As long as a different session does not have a write lock for an identifier, the session can acquire any number of read or write locks. Each lock request for the identifier acquires a new lock. The following statements acquire three write locks with the same identifier, then three read locks for the same identifier:

```
SELECT service_get_write_locks('ns', 'lock1', 'lock1', 'lock1', 0);
SELECT service_get_read_locks('ns', 'lock1', 'lock1', 'lock1', 0);
```

If you examine the Performance Schema `metadata_locks` table at this point, you will find that the session holds six distinct locks with the same `(ns, lock1)` identifier. (For details, see Locking Service Monitoring.)

Because the session holds at least one write lock on `(ns, lock1)`, no other session can acquire a lock for it, either read or write. If the session held only read locks for the identifier, other sessions could acquire read locks for it, but not write locks.

Locks for a single lock-acquisition call are acquired atomically, but atomicity does not hold across calls. Thus, for a statement such as the following, where `service_get_write_locks()` is called once per row of the result set, atomicity holds for each individual call, but not for the statement as a whole:

```
SELECT service_get_write_locks('ns', 'lock1', 'lock2', 0) FROM t1 WHERE ... ;
```

Caution

Because the locking service returns a separate lock for each successful request for a given lock identifier, it is possible for a single statement to acquire a large number of locks. For example:

```
INSERT INTO ... SELECT service_get_write_locks('ns', t1.col_name, 0) FROM t1;
```

These types of statements may have certain adverse effects. For example, if the statement fails part way through and rolls back, locks acquired up to the point of failure will still exist. If the intent is for there to be a correspondence between rows inserted and locks acquired, that intent will not be satisfied. Also, if it is important that locks are granted in a certain order, be aware that result set order may differ depending on which execution plan the optimizer chooses. For these reasons, it may be best to limit applications to a single lock-acquisition call per statement.

Locking Service Monitoring

The locking service is implemented using the MySQL Server metadata locks framework, so you monitor locking service locks acquired or waited for by examining the Performance Schema `metadata_locks` table.

First, enable the metadata lock instrument:

```
mysql> UPDATE performance_schema.setup_instruments SET ENABLED = 'YES'
    -> WHERE NAME = 'wait/lock/metadata/sql/mdl';
```

Then acquire some locks and check the contents of the `metadata_locks` table:

```
mysql> SELECT service_get_write_locks('mynamespace', 'lock1', 0);
+-----------------------------------------------------+
| service_get_write_locks('mynamespace', 'lock1', 0) |
+-----------------------------------------------------+
|                                                   1 |
+-----------------------------------------------------+
mysql> SELECT service_get_read_locks('mynamespace', 'lock2', 0);
+----------------------------------------------------+
| service_get_read_locks('mynamespace', 'lock2', 0) |
+----------------------------------------------------+
|                                                  1 |
+----------------------------------------------------+
mysql> SELECT OBJECT_TYPE, OBJECT_SCHEMA, OBJECT_NAME, LOCK_TYPE, LOCK_STATUS
    -> FROM performance_schema.metadata_locks
    -> WHERE OBJECT_TYPE = 'LOCKING SERVICE'\G
*************************** 1. row ***************************
  OBJECT_TYPE: LOCKING SERVICE
OBJECT_SCHEMA: mynamespace
  OBJECT_NAME: lock1
    LOCK_TYPE: EXCLUSIVE
  LOCK_STATUS: GRANTED
*************************** 2. row ***************************
  OBJECT_TYPE: LOCKING SERVICE
OBJECT_SCHEMA: mynamespace
  OBJECT_NAME: lock2
    LOCK_TYPE: SHARED
  LOCK_STATUS: GRANTED
```

Locking service locks have an `OBJECT_TYPE` value of `LOCKING SERVICE`. This is distinct from, for example, locks acquired with the `GET_LOCK()` function, which have an `OBJECT_TYPE` of `USER LEVEL LOCK`.

The lock namespace, name, and mode appear in the `OBJECT_SCHEMA`, `OBJECT_NAME`, and `LOCK_TYPE` columns. Read and write locks have `LOCK_TYPE` values of `SHARED` and `EXCLUSIVE`, respectively.

The `LOCK_STATUS` value is `GRANTED` for an acquired lock, `PENDING` for a lock that is being waited for. You will see `PENDING` if one session holds a write lock and another session is attempting to acquire a lock having the same identifier.

Locking Service UDF Interface Reference

The SQL interface to the locking service implements the user-defined functions described in this section. For usage examples, see Using the UDF Locking Interface.

The functions share these characteristics:

- The return value is nonzero for success. Otherwise, an error occurs.
- Namespace and lock names must be non-`NULL`, nonempty, and have a maximum length of 64 characters.

- Timeout values must be integers indicating how many seconds to wait to acquire locks before giving up with an error. If the timeout is 0, there is no waiting and the function produces an error if locks cannot be acquired immediately.

These locking service UDFs are available:

- `service_get_read_locks(namespace, lock_name[, lock_name] ..., timeout)`

  Acquires one or more read (shared) locks in the given namespace using the given lock names, timing out with an error if the locks are not acquired within the given timeout value.

- `service_get_write_locks(namespace, lock_name[, lock_name] ..., timeout)`

  Acquires one or more write (exclusive) locks in the given namespace using the given lock names, timing out with an error if the locks are not acquired within the given timeout value.

- `service_release_locks(namespace)`

  For the given namespace, releases all locks that were acquired within the current session using `service_get_read_locks()` and `service_get_write_locks()`.

  It is not an error for there to be no locks in the namespace.

28.3.2 The Keyring Service

MySQL Server supports a keyring service that enables internal server components and plugins to securely store sensitive information for later retrieval. This section describes how to use the keyring service functions to store, retrieve, and remove keys in the MySQL keyring keystore. An SQL interface to the keyring service functions is also available as a set of user-defined functions (UDFs); see Section 6.5.4.4, “General-Purpose Keyring Key-Management Functions”. For general keyring information, see Section 6.5.4, “The MySQL Keyring”.

The keyring service uses whatever underlying keyring plugin is enabled, if any. If no keyring plugin is enabled, keyring service calls fail.

A “record” in the keystore consists of data (the key itself) and a unique identifier through which the key is accessed. The identifier has two parts:

- `key_id`: The key ID or name. `key_id` values that begin with `mysql_` are reserved by MySQL Server.
- `user_id`: The session effective user ID. If there is no user context, this value can be `NULL`. The value need not actually be a “user”; the meaning depends on the application.

  Functions that implement the keyring UDF interface pass the value of `CURRENT_USER()` as the `user_id` value to keyring service functions.

The keyring service functions have these characteristics in common:

- Each function returns 0 for success, 1 for failure.
- The `key_id` and `user_id` arguments form a unique combination indicating which key in the keyring to use.
- The `key_type` argument provides additional information about the key, such as its encryption method or intended use.
- Keyring service functions treat key IDs, user names, types, and values as binary strings, so comparisons are case sensitive. For example, IDs of `MyKey` and `mykey` refer to different keys.

These keyring service functions are available:

- `my_key_fetch()`

Deobfuscates and retrieves a key from the keyring, along with its type. The function allocates the memory for the buffers used to store the returned key and key type. The caller should zero or obfuscate the memory when it is no longer needed, then free it.

Syntax:

```
bool my_key_fetch(const char *key_id, const char **key_type,
                  const char* user_id, void **key, size_t *key_len)
```

Arguments:

- `key_id`, `user_id`: Null-terminated strings that as a pair form a unique identifier indicating which key to fetch.
- `key_type`: The address of a buffer pointer. The function stores into it a pointer to a null-terminated string that provides additional information about the key (stored when the key was added).
- `key`: The address of a buffer pointer. The function stores into it a pointer to the buffer containing the fetched key data.
- `key_len`: The address of a variable into which the function stores the size in bytes of the `*key` buffer.

Return values:

Returns 0 for success, 1 for failure.

- `my_key_generate()`

  Generates a new random key of a given type and length and stores it in the keyring. The key has a length of `key_len` and is associated with the identifier formed from `key_id` and `user_id`. The type and length values must be consistent with the values supported by the underlying keyring plugin. See Section 6.5.4.3, “Supported Keyring Key Types”.

  Syntax:

  ```
  bool my_key_generate(const char *key_id, const char *key_type,
                       const char *user_id, size_t key_len)
  ```

  Arguments:

  - `key_id`, `user_id`: Null-terminated strings that as a pair form a unique identifier for the key to be generated.
  - `key_type`: A null-terminated string that provides additional information about the key.
  - `key_len`: The size in bytes of the key to be generated.

  Return values:

  Returns 0 for success, 1 for failure.

- `my_key_remove()`

  Removes a key from the keyring.

  Syntax:

  ```
  bool my_key_remove(const char *key_id, const char* user_id)
  ```

Arguments:

- `key_id`, `user_id`: Null-terminated strings that as a pair form a unique identifier for the key to be removed.

Return values:

Returns 0 for success, 1 for failure.

- `my_key_store()`

  Obfuscates and stores a key in the keyring.

  Syntax:

```
bool my_key_store(const char *key_id, const char *key_type,
                  const char* user_id, void *key, size_t key_len)
```

  Arguments:

  - `key_id`, `user_id`: Null-terminated strings that as a pair form a unique identifier for the key to be stored.
  - `key_type`: A null-terminated string that provides additional information about the key.
  - `key`: The buffer containing the key data to be stored.
  - `key_len`: The size in bytes of the `key` buffer.

  Return values:

  Returns 0 for success, 1 for failure.

28.4 Adding New Functions to MySQL

There are three ways to add new functions to MySQL:

- You can add functions through the user-defined function (UDF) interface. User-defined functions are compiled as library files and then added to and removed from the server dynamically using the `CREATE FUNCTION` and `DROP FUNCTION` statements. See Section 13.7.3.1, “CREATE FUNCTION Syntax for User-Defined Functions”.
- You can add functions as native (built-in) MySQL functions. Native functions are compiled into the `mysqld` server and become available on a permanent basis.
- Another way to add functions is by creating stored functions. These are written using SQL statements rather than by compiling object code. The syntax for writing stored functions is not covered here. See Section 23.2, “Using Stored Routines (Procedures and Functions)”.

Each method of creating compiled functions has advantages and disadvantages:

- If you write user-defined functions, you must install object files in addition to the server itself. If you compile your function into the server, you need not do that.
- Native functions require you to modify a source distribution. UDFs do not. You can add UDFs to a binary MySQL distribution. No access to MySQL source is necessary.
- If you upgrade your MySQL distribution, you can continue to use your previously installed UDFs, unless you upgrade to a newer version for which the UDF interface changes. For native functions, you must repeat your modifications each time you upgrade.

Whichever method you use to add new functions, they can be invoked in SQL statements just like native functions such as `ABS()` or `SOUNDEX()`.

See Section 9.2.4, “Function Name Parsing and Resolution”, for the rules describing how the server interprets references to different kinds of functions.

The following sections describe features of the UDF interface, provide instructions for writing UDFs, discuss security precautions that MySQL takes to prevent UDF misuse, and describe how to add native MySQL functions.

For example source code that illustrates how to write UDFs, take a look at the `sql/udf_example.cc` file that is provided in MySQL source distributions.

Note

The MySQL source code contains internal documentation written using Doxygen. This documentation is useful for understanding how MySQL works from a developer perspective. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, “Generating MySQL Doxygen Documentation Content”.

28.4.1 Features of the User-Defined Function Interface

The MySQL interface for user-defined functions provides the following features and capabilities:

- Functions can return string, integer, or real values and can accept arguments of those same types.
- You can define simple functions that operate on a single row at a time, or aggregate functions that operate on groups of rows.
- Information is provided to functions that enables them to check the number, types, and names of the arguments passed to them.
- You can tell MySQL to coerce arguments to a given type before passing them to a function.
- You can indicate that a function returns `NULL` or that an error occurred.

28.4.2 Adding a New User-Defined Function

For the UDF mechanism to work, functions must be written in C or C++ and your operating system must support dynamic loading. MySQL source distributions include a file `sql/udf_example.cc` that defines five UDF functions. Consult this file to see how UDF calling conventions work. The `include/mysql_com.h` header file defines UDF-related symbols and data structures, although you need not include this header file directly; it is included by `mysql.h`.

A UDF contains code that becomes part of the running server, so when you write a UDF, you are bound by any and all constraints that apply to writing server code. For example, you may have problems if you attempt to use functions from the `libstdc++` library. These constraints may change in future versions of the server, so it is possible that server upgrades will require revisions to UDFs that were originally written for older servers. For information about these constraints, see Section 2.8.4, “MySQL Source-Configuration Options”, and Section 2.8.5, “Dealing with Problems Compiling MySQL”.

To be able to use UDFs, you must link `mysqld` dynamically. If you want to use a UDF that needs to access symbols from `mysqld` (for example, the `metaphone` function in `sql/udf_example.cc` uses `default_charset_info`), you must link the program with `-rdynamic` (see `man dlopen`).

For each function that you want to use in SQL statements, you should define corresponding C (or C ++) functions. In the following discussion, the name “xxx” is used for an example function name. To

distinguish between SQL and C/C++ usage, `XXX()` (uppercase) indicates an SQL function call, and `xxx()` (lowercase) indicates a C/C++ function call.

Note

When using C++ you can encapsulate your C functions within:

```
extern "C" { ... }
```

This ensures that your C++ function names remain readable in the completed UDF.

The following list describes the C/C++ functions that you write to implement the interface for a function named `XXX()`. The main function, `xxx()`, is required. In addition, a UDF requires at least one of the other functions described here, for reasons discussed in Section 28.4.2.6, “UDF Security Precautions”.

- `xxx()`

  The main function. This is where the function result is computed. The correspondence between the SQL function data type and the return type of your C/C++ function is shown here.

SQL Type	C/C++ Type
`STRING`	`char *`
`INTEGER`	`long long`
`REAL`	`double`

  It is also possible to declare a `DECIMAL` function, but currently the value is returned as a string, so you should write the UDF as though it were a `STRING` function. `ROW` functions are not implemented.

- `xxx_init()`

  The initialization function for `xxx()`. If present, it can be used for the following purposes:

  - To check the number of arguments to `XXX()`.
  - To verify that the arguments are of a required type or, alternatively, to tell MySQL to coerce arguments to the required types when the main function is called.
  - To allocate any memory required by the main function.
  - To specify the maximum length of the result.
  - To specify (for `REAL` functions) the maximum number of decimal places in the result.
  - To specify whether the result can be `NULL`.

- `xxx_deinit()`

  The deinitialization function for `xxx()`. If present, it should deallocate any memory allocated by the initialization function.

When an SQL statement invokes `XXX()`, MySQL calls the initialization function `xxx_init()` to let it perform any required setup, such as argument checking or memory allocation. If `xxx_init()` returns an error, MySQL aborts the SQL statement with an error message and does not call the main or deinitialization functions. Otherwise, MySQL calls the main function `xxx()` once for each row. After all rows have been processed, MySQL calls the deinitialization function `xxx_deinit()` so that it can perform any required cleanup.

For aggregate functions that work like `SUM()`, you must also provide the following functions:

- `xxx_clear()`

  Reset the current aggregate value but do not insert the argument as the initial aggregate value for a new group.

- `xxx_add()`

  Add the argument to the current aggregate value.

MySQL handles aggregate UDFs as follows:

1. Call `xxx_init()` to let the aggregate function allocate any memory it needs for storing results.
2. Sort the table according to the `GROUP BY` expression.
3. Call `xxx_clear()` for the first row in each new group.
4. Call `xxx_add()` for each row that belongs in the same group.
5. Call `xxx()` to get the result for the aggregate when the group changes or after the last row has been processed.
6. Repeat steps 3 to 5 until all rows has been processed
7. Call `xxx_deinit()` to let the UDF free any memory it has allocated.

All functions must be thread-safe. This includes not just the main function, but the initialization and deinitialization functions as well, and also the additional functions required by aggregate functions. A consequence of this requirement is that you are not permitted to allocate any global or static variables that change! If you need memory, you should allocate it in `xxx_init()` and free it in `xxx_deinit()`.

28.4.2.1 UDF Calling Sequences for Simple Functions

This section describes the different functions that you need to define when you create a simple UDF. Section 28.4.2, “Adding a New User-Defined Function”, describes the order in which MySQL calls these functions.

The main `xxx()` function should be declared as shown in this section. Note that the return type and parameters differ, depending on whether you declare the SQL function `XXX()` to return `STRING`, `INTEGER`, or `REAL` in the `CREATE FUNCTION` statement:

For `STRING` functions:

```
char *xxx(UDF_INIT *initid, UDF_ARGS *args,
          char *result, unsigned long *length,
          char *is_null, char *error);
```

For `INTEGER` functions:

```
long long xxx(UDF_INIT *initid, UDF_ARGS *args,
              char *is_null, char *error);
```

For `REAL` functions:

```
double xxx(UDF_INIT *initid, UDF_ARGS *args,
              char *is_null, char *error);
```

`DECIMAL` functions return string values and should be declared the same way as `STRING` functions. `ROW` functions are not implemented.

The initialization and deinitialization functions are declared like this:

```
bool xxx_init(UDF_INIT *initid, UDF_ARGS *args, char *message);

void xxx_deinit(UDF_INIT *initid);
```

The `initid` parameter is passed to all three functions. It points to a `UDF_INIT` structure that is used to communicate information between functions. The `UDF_INIT` structure members follow. The initialization function should fill in any members that it wishes to change. (To use the default for a member, leave it unchanged.)

- `bool maybe_null`

  `xxx_init()` should set `maybe_null` to `1` if `xxx()` can return `NULL`. The default value is `1` if any of the arguments are declared `maybe_null`.

- `unsigned int decimals`

  The number of decimal digits to the right of the decimal point. The default value is the maximum number of decimal digits in the arguments passed to the main function. For example, if the function is passed `1.34`, `1.345`, and `1.3`, the default would be 3, because `1.345` has 3 decimal digits.

  For arguments that have no fixed number of decimals, the `decimals` value is set to 31, which is 1 more than the maximum number of decimals permitted for the `DECIMAL`, `FLOAT`, and `DOUBLE` data types. This value is available as the constant `NOT_FIXED_DEC` in the `mysql_com.h` header file.

  A `decimals` value of 31 is used for arguments in cases such as a `FLOAT` or `DOUBLE` column declared without an explicit number of decimals (for example, `FLOAT` rather than `FLOAT(10,3)`) and for floating-point constants such as `1345E-3`. It is also used for string and other nonnumber arguments that might be converted within the function to numeric form.

  The value to which the `decimals` member is initialized is only a default. It can be changed within the function to reflect the actual calculation performed. The default is determined such that the largest number of decimals of the arguments is used. If the number of decimals is `NOT_FIXED_DEC` for even one of the arguments, that is the value used for `decimals`.

- `unsigned int max_length`

  The maximum length of the result. The default `max_length` value differs depending on the result type of the function. For string functions, the default is the length of the longest argument. For integer functions, the default is 21 digits. For real functions, the default is 13 plus the number of decimal digits indicated by `initid->decimals`. (For numeric functions, the length includes any sign or decimal point characters.)

  If you want to return a blob value, you can set `max_length` to 65KB or 16MB. This memory is not allocated, but the value is used to decide which data type to use if there is a need to temporarily store the data.

- `char *ptr`

  A pointer that the function can use for its own purposes. For example, functions can use `initid->ptr` to communicate allocated memory among themselves. `xxx_init()` should allocate the memory and assign it to this pointer:

  ```
  initid->ptr = allocated_memory;
  ```

  In `xxx()` and `xxx_deinit()`, refer to `initid->ptr` to use or deallocate the memory.

- `bool const_item`

  `xxx_init()` should set `const_item` to `1` if `xxx()` always returns the same value and to `0` otherwise.

28.4.2.2 UDF Calling Sequences for Aggregate Functions

This section describes the different functions that you need to define when you create an aggregate UDF. Section 28.4.2, “Adding a New User-Defined Function”, describes the order in which MySQL calls these functions.

- `xxx_reset()`

  This function is called when MySQL finds the first row in a new group. It should reset any internal summary variables and then use the given `UDF_ARGS` argument as the first value in your internal summary value for the group. Declare `xxx_reset()` as follows:

  ```
  void xxx_reset(UDF_INIT *initid, UDF_ARGS *args,
                 char *is_null, char *error);
  ```

  `xxx_reset()` is not needed or used in MySQL 8.0, in which the UDF interface uses `xxx_clear()` instead. However, you can define both `xxx_reset()` and `xxx_clear()` if you want to have your UDF work with older versions of the server. (If you do include both functions, the `xxx_reset()` function in many cases can be implemented internally by calling `xxx_clear()` to reset all variables, and then calling `xxx_add()` to add the `UDF_ARGS` argument as the first value in the group.)

- `xxx_clear()`

  This function is called when MySQL needs to reset the summary results. It is called at the beginning for each new group but can also be called to reset the values for a query where there were no matching rows. Declare `xxx_clear()` as follows:

  ```
  void xxx_clear(UDF_INIT *initid, char *is_null, char *error);
  ```

  `is_null` is set to point to `CHAR(0)` before calling `xxx_clear()`.

  If something went wrong, you can store a value in the variable to which the `error` argument points. `error` points to a single-byte variable, not to a string buffer.

  `xxx_clear()` is required by MySQL 8.0.

- `xxx_add()`

  This function is called for all rows that belong to the same group. You should use it to add the value in the `UDF_ARGS` argument to your internal summary variable.

  ```
  void xxx_add(UDF_INIT *initid, UDF_ARGS *args,
               char *is_null, char *error);
  ```

The `xxx()` function for an aggregate UDF should be declared the same way as for a nonaggregate UDF. See Section 28.4.2.1, “UDF Calling Sequences for Simple Functions”.

For an aggregate UDF, MySQL calls the `xxx()` function after all rows in the group have been processed. You should normally never access its `UDF_ARGS` argument here but instead return a value based on your internal summary variables.

Return value handling in `xxx()` should be done the same way as for a nonaggregate UDF. See Section 28.4.2.4, “UDF Return Values and Error Handling”.

The `xxx_reset()` and `xxx_add()` functions handle their `UDF_ARGS` argument the same way as functions for nonaggregate UDFs. See Section 28.4.2.3, “UDF Argument Processing”.

The pointer arguments to `is_null` and `error` are the same for all calls to `xxx_reset()`, `xxx_clear()`, `xxx_add()` and `xxx()`. You can use this to remember that you got an error or whether the `xxx()` function should return `NULL`. You should not store a string into `*error`! `error` points to a single-byte variable, not to a string buffer.

`*is_null` is reset for each group (before calling `xxx_clear()`). `*error` is never reset.

If `*is_null` or `*error` are set when `xxx()` returns, MySQL returns `NULL` as the result for the group function.

28.4.2.3 UDF Argument Processing

The `args` parameter points to a `UDF_ARGS` structure that has the members listed here:

- `unsigned int arg_count`

  The number of arguments. Check this value in the initialization function if you require your function to be called with a particular number of arguments. For example:

```
if (args->arg_count != 2)
{
    strcpy(message,"XXX() requires two arguments");
    return 1;
}
```

  For other `UDF_ARGS` member values that are arrays, array references are zero-based. That is, refer to array members using index values from 0 to `args->arg_count` − 1.

- `enum Item_result *arg_type`

  A pointer to an array containing the types for each argument. The possible type values are `STRING_RESULT`, `INT_RESULT`, `REAL_RESULT`, and `DECIMAL_RESULT`.

  To make sure that arguments are of a given type and return an error if they are not, check the `arg_type` array in the initialization function. For example:

```
if (args->arg_type[0] != STRING_RESULT ||
    args->arg_type[1] != INT_RESULT)
{
    strcpy(message,"XXX() requires a string and an integer");
    return 1;
}
```

  Arguments of type `DECIMAL_RESULT` are passed as strings, so you should handle them the same way as `STRING_RESULT` values.

  As an alternative to requiring your function's arguments to be of particular types, you can use the initialization function to set the `arg_type` elements to the types you want. This causes MySQL to coerce arguments to those types for each call to `xxx()`. For example, to specify that the first two arguments should be coerced to string and integer, respectively, do this in `xxx_init()`:

```
args->arg_type[0] = STRING_RESULT;
args->arg_type[1] = INT_RESULT;
```

  Exact-value decimal arguments such as `1.3` or `DECIMAL` column values are passed with a type of `DECIMAL_RESULT`. However, the values are passed as strings. If you want to receive a number, use the initialization function to specify that the argument should be coerced to a `REAL_RESULT` value:

```
args->arg_type[2] = REAL_RESULT;
```

- `char **args`

  `args->args` communicates information to the initialization function about the general nature of the arguments passed to your function. For a constant argument `i`, `args->args[i]` points to the argument value. (See later for instructions on how to access the value properly.) For a nonconstant argument, `args->args[i]` is `0`. A constant argument is an expression that uses only constants,

such as `3` or `4*7-2` or `SIN(3.14)`. A nonconstant argument is an expression that refers to values that may change from row to row, such as column names or functions that are called with nonconstant arguments.

For each invocation of the main function, `args->args` contains the actual arguments that are passed for the row currently being processed.

If argument `i` represents `NULL`, `args->args[i]` is a null pointer (0). If the argument is not `NULL`, functions can refer to it as follows:

- An argument of type `STRING_RESULT` is given as a string pointer plus a length, to enable handling of binary data or data of arbitrary length. The string contents are available as `args->args[i]` and the string length is `args->lengths[i]`. Do not assume that the string is null-terminated.
- For an argument of type `INT_RESULT`, you must cast `args->args[i]` to a `long long` value:

```
long long int_val;
int_val = *((long long*) args->args[i]);
```

- For an argument of type `REAL_RESULT`, you must cast `args->args[i]` to a `double` value:

```
double    real_val;
real_val = *((double*) args->args[i]);
```

- For an argument of type `DECIMAL_RESULT`, the value is passed as a string and should be handled like a `STRING_RESULT` value.
- `ROW_RESULT` arguments are not implemented.

- `unsigned long *lengths`

  For the initialization function, the `lengths` array indicates the maximum string length for each argument. You should not change these. For each invocation of the main function, `lengths` contains the actual lengths of any string arguments that are passed for the row currently being processed. For arguments of types `INT_RESULT` or `REAL_RESULT`, `lengths` still contains the maximum length of the argument (as for the initialization function).

- `char *maybe_null`

  For the initialization function, the `maybe_null` array indicates for each argument whether the argument value might be null (0 if no, 1 if yes).

- `char **attributes`

  `args->attributes` communicates information about the names of the UDF arguments. For argument `i`, the attribute name is available as a string in `args->attributes[i]` and the attribute length is `args->attribute_lengths[i]`. Do not assume that the string is null-terminated.

  By default, the name of a UDF argument is the text of the expression used to specify the argument. For UDFs, an argument may also have an optional `[AS] alias_name` clause, in which case the argument name is *`alias_name`*. The `attributes` value for each argument thus depends on whether an alias was given.

  Suppose that a UDF `my_udf()` is invoked as follows:

```
SELECT my_udf(expr1, expr2 AS alias1, expr3 alias2);
```

  In this case, the `attributes` and `attribute_lengths` arrays will have these values:

```
args->attributes[0] = "expr1"
```

```
args->attribute_lengths[0] = 5

args->attributes[1] = "alias1"
args->attribute_lengths[1] = 6

args->attributes[2] = "alias2"
args->attribute_lengths[2] = 6
```

- `unsigned long *attribute_lengths`

  The `attribute_lengths` array indicates the length of each argument name.

28.4.2.4 UDF Return Values and Error Handling

The initialization function should return `0` if no error occurred and `1` otherwise. If an error occurs, `xxx_init()` should store a null-terminated error message in the `message` parameter. The message is returned to the client. The message buffer is `MYSQL_ERRMSG_SIZE` characters long, but you should try to keep the message to less than 80 characters so that it fits the width of a standard terminal screen.

The return value of the main function `xxx()` is the function value, for `long long` and `double` functions. A string function should return a pointer to the result and set `*length` to the length (in bytes) of the return value. For example:

```
memcpy(result, "result string", 13);
*length = 13;
```

MySQL passes a buffer to the `xxx()` function using the `result` parameter. This buffer is sufficiently long to hold 255 characters, which can be multibyte characters. The `xxx()` function can store the result in this buffer if it fits, in which case the return value should be a pointer to the buffer. If the function stores the result in a different buffer, it should return a pointer to that buffer.

If your string function does not use the supplied buffer (for example, if it needs to return a string longer than 255 characters), you must allocate the space for your own buffer with `malloc()` in your `xxx_init()` function or your `xxx()` function and free it in your `xxx_deinit()` function. You can store the allocated memory in the `ptr` slot in the `UDF_INIT` structure for reuse by future `xxx()` calls. See Section 28.4.2.1, “UDF Calling Sequences for Simple Functions”.

To indicate a return value of `NULL` in the main function, set `*is_null` to `1`:

```
*is_null = 1;
```

To indicate an error return in the main function, set `*error` to `1`:

```
*error = 1;
```

If `xxx()` sets `*error` to `1` for any row, the function value is `NULL` for the current row and for any subsequent rows processed by the statement in which `XXX()` was invoked. (`xxx()` is not even called for subsequent rows.)

28.4.2.5 UDF Compiling and Installing

Files implementing UDFs must be compiled and installed on the host where the server runs. This process is described below for the example UDF file `sql/udf_example.cc` that is included in MySQL source distributions.

If a UDF will be referred to in statements that will be replicated to slave servers, you must ensure that every slave also has the function available. Otherwise, replication will fail on the slaves when they attempt to invoke the function.

The immediately following instructions are for Unix. Instructions for Windows are given later in this section.

The `udf_example.cc` file contains the following functions:

- `metaphon()` returns a metaphon string of the string argument. This is something like a soundex string, but it is more tuned for English.
- `myfunc_double()` returns the sum of the ASCII values of the characters in its arguments, divided by the sum of the length of its arguments.
- `myfunc_int()` returns the sum of the length of its arguments.
- `sequence([const int])` returns a sequence starting from the given number or 1 if no number has been given.
- `lookup()` returns the IP address for a host name.
- `reverse_lookup()` returns the host name for an IP address. The function may be called either with a single string argument of the form `'xxx.xxx.xxx.xxx'` or with four numbers.
- `avgcost()` returns an average cost. This is an aggregate function.

A dynamically loadable file should be compiled as a sharable library file, using a command something like this:

```
shell> gcc -shared -o udf_example.so udf_example.cc
```

If you are using `gcc` with `CMake` (which is how MySQL is configured), you should be able to create `udf_example.so` with a simpler command:

```
shell> make udf_example
```

After you compile a shared object containing UDFs, you must install it and tell MySQL about it. Compiling a shared object from `udf_example.cc` using `gcc` directly produces a file named `udf_example.so`. Copy the shared object to the server's plugin directory and name it `udf_example.so`. This directory is given by the value of the `plugin_dir` system variable.

On some systems, the `ldconfig` program that configures the dynamic linker does not recognize a shared object unless its name begins with `lib`. In this case you should rename a file such as `udf_example.so` to `libudf_example.so`.

On Windows, you can compile user-defined functions by using the following procedure:

1. Obtain a MySQL source distribution. See Section 2.1.2, “How to Get MySQL”.
2. Obtain the `CMake` build utility, if necessary, from http://www.cmake.org. (Version 2.6 or later is required).
3. In the source tree, look in the `sql` directory. There are files named `udf_example.def` `udf_example.cc` there. Copy both files from this directory to your working directory.
4. Create a `CMake makefile` (`CMakeLists.txt`) with these contents:

   ```
   PROJECT(udf_example)

   # Path for MySQL include directory
   INCLUDE_DIRECTORIES("c:/mysql/include")

   ADD_DEFINITIONS("-DHAVE_DLOPEN")
   ADD_LIBRARY(udf_example MODULE udf_example.cc udf_example.def)
   TARGET_LINK_LIBRARIES(udf_example wsock32)
   ```

5. Create the VC project and solution files:

```
cmake -G "<Generator>"
```

Invoking `cmake --help` shows you a list of valid Generators.

6. Create `udf_example.dll`:

```
devenv udf_example.sln /build Release
```

After the shared library file has been installed, notify `mysqld` about the new functions with the following statements. If library files have a suffix different from `.so` on your system, substitute the correct suffix throughout (for example, `.dll` on Windows).

```
mysql> CREATE FUNCTION metaphon RETURNS STRING SONAME 'udf_example.so';
mysql> CREATE FUNCTION myfunc_double RETURNS REAL SONAME 'udf_example.so';
mysql> CREATE FUNCTION myfunc_int RETURNS INTEGER SONAME 'udf_example.so';
mysql> CREATE FUNCTION sequence RETURNS INTEGER SONAME 'udf_example.so';
mysql> CREATE FUNCTION lookup RETURNS STRING SONAME 'udf_example.so';
mysql> CREATE FUNCTION reverse_lookup
    ->        RETURNS STRING SONAME 'udf_example.so';
mysql> CREATE AGGREGATE FUNCTION avgcost
    ->        RETURNS REAL SONAME 'udf_example.so';
```

Once installed, a function remains installed until it is uninstalled.

To delete functions, use `DROP FUNCTION`:

```
mysql> DROP FUNCTION metaphon;
mysql> DROP FUNCTION myfunc_double;
mysql> DROP FUNCTION myfunc_int;
mysql> DROP FUNCTION sequence;
mysql> DROP FUNCTION lookup;
mysql> DROP FUNCTION reverse_lookup;
mysql> DROP FUNCTION avgcost;
```

The `CREATE FUNCTION` and `DROP FUNCTION` statements update the `func` system table in the `mysql` database. The function's name, type and shared library name are saved in the table. You must have the `INSERT` or `DELETE` privilege for the `mysql` database to create or drop functions, respectively.

You should not use `CREATE FUNCTION` to add a function that has previously been created. If you need to reinstall a function, you should remove it with `DROP FUNCTION` and then reinstall it with `CREATE FUNCTION`. You would need to do this, for example, if you recompile a new version of your function, so that `mysqld` gets the new version. Otherwise, the server continues to use the old version.

An active function is one that has been loaded with `CREATE FUNCTION` and not removed with `DROP FUNCTION`. All active functions are reloaded each time the server starts, unless you start `mysqld` with the `--skip-grant-tables` option. In this case, UDF initialization is skipped and UDFs are unavailable.

28.4.2.6 UDF Security Precautions

MySQL takes several measures to prevent misuse of user-defined functions.

UDF library files cannot be placed in arbitrary directories. They must be located in the server's plugin directory. This directory is given by the value of the `plugin_dir` system variable.

To use `CREATE FUNCTION` or `DROP FUNCTION`, you must have the `INSERT` or `DELETE` privilege, respectively, for the `mysql` database. This is necessary because those statements add and delete rows from the `mysql.func` table.

UDFs should have at least one symbol defined in addition to the `xxx` symbol that corresponds to the main `xxx()` function. These auxiliary symbols correspond to the `xxx_init()`, `xxx_deinit()`,

`xxx_reset()`, `xxx_clear()`, and `xxx_add()` functions. `mysqld` also supports an `--allow-suspicious-udfs` option that controls whether UDFs that have only an `xxx` symbol can be loaded. By default, the option is off, to prevent attempts at loading functions from shared library files other than those containing legitimate UDFs. If you have older UDFs that contain only the `xxx` symbol and that cannot be recompiled to include an auxiliary symbol, it may be necessary to specify the `--allow-suspicious-udfs` option. Otherwise, you should avoid enabling this capability.

28.4.3 Adding a New Native Function

To add a new native MySQL function, use the procedure described here, which requires that you use a source distribution. You cannot add native functions to a binary distribution because it is necessary to modify MySQL source code and compile MySQL from the modified source. If you migrate to another version of MySQL (for example, when a new version is released), you must repeat the procedure with the new version.

If the new native function will be referred to in statements that will be replicated to slave servers, you must ensure that every slave server also has the function available. Otherwise, replication will fail on the slaves when they attempt to invoke the function.

To add a new native function, follow these steps to modify source files in the `sql` directory:

1. Create a subclass for the function in `item_create.cc`:

   - If the function takes a fixed number of arguments, create a subclass of `Create_func_arg0`, `Create_func_arg1`, `Create_func_arg2`, or `Create_func_arg3`, respectively, depending on whether the function takes zero, one, two, or three arguments. For examples, see the `Create_func_uuid`, `Create_func_abs`, `Create_func_pow`, and `Create_func_lpad` classes.

   - If the function takes a variable number of arguments, create a subclass of `Create_native_func`. For an example, see `Create_func_concat`.

2. To provide a name by which the function can be referred to in SQL statements, register the name in `item_create.cc` by adding a line to this array:

   ```
   static Native_func_registry func_array[]
   ```

   You can register several names for the same function. For example, see the lines for `"LCASE"` and `"LOWER"`, which are aliases for `Create_func_lcase`.

3. In `item_func.h`, declare a class inheriting from `Item_num_func` or `Item_str_func`, depending on whether your function returns a number or a string.

4. In `item_func.cc`, add one of the following declarations, depending on whether you are defining a numeric or string function:

   ```
   double   Item_func_newname::val()
   longlong Item_func_newname::val_int()
   String  *Item_func_newname::Str(String *str)
   ```

   If you inherit your object from any of the standard items (like `Item_num_func`), you probably only have to define one of these functions and let the parent object take care of the other functions. For example, the `Item_str_func` class defines a `val()` function that executes `atof()` on the value returned by `::str()`.

5. If the function is nondeterministic, include the following statement in the item constructor to indicate that function results should not be cached:

   ```
   current_thd->lex->safe_to_cache_query=0;
   ```

A function is nondeterministic if, given fixed values for its arguments, it can return different results for different invocations.

6. You should probably also define the following object function:

```
void Item_func_newname::fix_length_and_dec()
```

This function should at least calculate `max_length` based on the given arguments. `max_length` is the maximum number of characters the function may return. This function should also set `maybe_null = 0` if the main function cannot return a `NULL` value. The function can check whether any of the function arguments can return `NULL` by checking the arguments' `maybe_null` variable. Look at `Item_func_mod::fix_length_and_dec` for a typical example of how to do this.

All functions must be thread-safe. In other words, do not use any global or static variables in the functions without protecting them with mutexes.

If you want to return `NULL` from `::val()`, `::val_int()`, or `::str()`, you should set `null_value` to 1 and return 0.

For `::str()` object functions, there are additional considerations to be aware of:

- The `String *str` argument provides a string buffer that may be used to hold the result. (For more information about the `String` type, take a look at the `sql_string.h` file.)
- The `::str()` function should return the string that holds the result, or `(char*) 0` if the result is `NULL`.
- All current string functions try to avoid allocating any memory unless absolutely necessary!

28.5 Debugging and Porting MySQL

This section helps you port MySQL to other operating systems. Do check the list of currently supported operating systems first. See http://www.mysql.com/support/supportedplatforms/database.html. If you have created a new port of MySQL, please let us know so that we can list it here and on our Web site (http://www.mysql.com/), recommending it to other users.

Note

If you create a new port of MySQL, you are free to copy and distribute it under the GPL license, but it does not make you a copyright holder of MySQL.

A working POSIX thread library is needed for the server.

To build MySQL from source, your system must satisfy the tool requirements listed at Section 2.8, “Installing MySQL from Source”.

If you run into problems with a new port, you may have to do some debugging of MySQL! See Section 28.5.1, “Debugging a MySQL Server”.

Note

Before you start debugging `mysqld`, first get the test program `mysys/thr_lock` to work. This ensures that your thread installation has even a remote chance to work!

Note

The MySQL source code contains internal documentation written using Doxygen. This documentation is useful for understanding how MySQL works

from a developer perspective. The generated Doxygen content is available at http://dev.mysql.com/doc/dev/mysql-server/latest/. It is also possible to generate this content locally from a MySQL source distribution using the instructions at Section 2.8.7, “Generating MySQL Doxygen Documentation Content”.

28.5.1 Debugging a MySQL Server

If you are using some functionality that is very new in MySQL, you can try to run `mysqld` with the `--skip-new` (which disables all new, potentially unsafe functionality). See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.

If `mysqld` does not want to start, verify that you have no `my.cnf` files that interfere with your setup! You can check your `my.cnf` arguments with `mysqld --print-defaults` and avoid using them by starting with `mysqld --no-defaults ...`.

If `mysqld` starts to eat up CPU or memory or if it “hangs,” you can use `mysqladmin processlist status` to find out if someone is executing a query that takes a long time. It may be a good idea to run `mysqladmin -i10 processlist status` in some window if you are experiencing performance problems or problems when new clients cannot connect.

The command `mysqladmin debug` dumps some information about locks in use, used memory and query usage to the MySQL log file. This may help solve some problems. This command also provides some useful information even if you have not compiled MySQL for debugging!

If the problem is that some tables are getting slower and slower you should try to optimize the table with `OPTIMIZE TABLE` or `myisamchk`. See Chapter 5, *MySQL Server Administration*. You should also check the slow queries with `EXPLAIN`.

You should also read the OS-specific section in this manual for problems that may be unique to your environment. See Section 2.1, “General Installation Guidance”.

28.5.1.1 Compiling MySQL for Debugging

If you have some very specific problem, you can always try to debug MySQL. To do this you must configure MySQL with the `-DWITH_DEBUG=1` option. You can check whether MySQL was compiled with debugging by doing: `mysqld --help`. If the `--debug` flag is listed with the options then you have debugging enabled. `mysqladmin ver` also lists the `mysqld` version as `mysql ... --debug` in this case.

If `mysqld` stops crashing when you configure it with the `-DWITH_DEBUG=1` CMake option, you probably have found a compiler bug or a timing bug within MySQL. In this case, you can try to add `-g` using the `CMAKE_C_FLAGS` and `CMAKE_CXX_FLAGS` CMake options and not use `-DWITH_DEBUG=1`. If `mysqld` dies, you can at least attach to it with `gdb` or use `gdb` on the core file to find out what happened.

When you configure MySQL for debugging you automatically enable a lot of extra safety check functions that monitor the health of `mysqld`. If they find something “unexpected,” an entry is written to `stderr`, which `mysqld_safe` directs to the error log! This also means that if you are having some unexpected problems with MySQL and are using a source distribution, the first thing you should do is to configure MySQL for debugging! (The second thing is to send mail to a MySQL mailing list and ask for help. See Section 1.6.2, “MySQL Mailing Lists”. If you believe that you have found a bug, please use the instructions at Section 1.7, “How to Report Bugs or Problems”.

In the Windows MySQL distribution, `mysqld.exe` is by default compiled with support for trace files.

28.5.1.2 Creating Trace Files

If the `mysqld` server does not start or it crashes easily, you can try to create a trace file to find the problem.

To do this, you must have a `mysqld` that has been compiled with debugging support. You can check this by executing `mysqld -V`. If the version number ends with `-debug`, it is compiled with support for trace files. (On Windows, the debugging server is named `mysqld-debug` rather than `mysqld`.)

Start the `mysqld` server with a trace log in `/tmp/mysqld.trace` on Unix or `\mysqld.trace` on Windows:

```
shell> mysqld --debug
```

On Windows, you should also use the `--standalone` flag to not start `mysqld` as a service. In a console window, use this command:

```
C:\> mysqld-debug --debug --standalone
```

After this, you can use the `mysql.exe` command-line tool in a second console window to reproduce the problem. You can stop the `mysqld` server with `mysqladmin shutdown`.

The trace file can become **very large**! To generate a smaller trace file, you can use debugging options something like this:

```
mysqld --debug=d,info,error,query,general,where:O,/tmp/mysqld.trace
```

This only prints information with the most interesting tags to the trace file.

If you make a bug report about this, please only send the lines from the trace file to the appropriate mailing list where something seems to go wrong! If you cannot locate the wrong place, you can open a bug report and upload the trace file to the report, so that a MySQL developer can take a look at it. For instructions, see Section 1.7, “How to Report Bugs or Problems”.

The trace file is made with the **DBUG** package by Fred Fish. See Section 28.5.3, “The DBUG Package”.

28.5.1.3 Using WER with PDB to create a Windows crashdump

Program Database files (extension `pdb`) are included in the **ZIP Archive Debug Binaries & Test Suite** distribution of MySQL. These files provide information for debugging your MySQL installation in the event of a problem. This is a separate download from the standard MSI or Zip file.

Note

The PDB files are available in a separate file labeled "ZIP Archive Debug Binaries & Test Suite".

The PDB file contains more detailed information about `mysqld` and other tools that enables more detailed trace and dump files to be created. You can use these with `WinDbg` or Visual Studio to debug `mysqld`.

Note

The older *Dr. Watson* debugging tool was removed in Microsoft Vista, with `WinDbg` being a common alternative.

For more information on PDB files, see Microsoft Knowledge Base Article 121366. For more information on the debugging options available, see Debugging Tools for Windows.

To use WinDbg, either install the full Windows Driver Kit (WDK) or install the standalone version.

Important

The `.exe` and `.pbd` files must be an exact match (both version number and MySQL server edition) or WinDBG will complain while attempting to load the symbols.

1. To generate a minidump `mysqld.dmp`, enable the `core-file` option under the [mysqld] section in `my.ini`. Restart the MySQL server after making these changes.

2. Create a directory to store the generated files, such as `c:\symbols`

3. Determine the path to your `windbg.exe` executable using the Find GUI or from the command line, for example: `dir /s /b windbg.exe` -- a common default is *C:\Program Files\Debugging Tools for Windows (x64)\windbg.exe*

4. Launch `windbg.exe` giving it the paths to `mysqld.exe`, `mysqld.pdb`, `mysqld.dmp`, and the source code. Alternatively, pass in each path from the WinDbg GUI. For example:

```
windbg.exe -i "C:\mysql-8.0.3-winx64\bin\"^
 -z "C:\mysql-8.0.3-winx64\data\mysqld.dmp"^
 -srcpath "E:\ade\mysql_archives\8.0\8.0.3\mysql-8.0.3"^
 -y "C:\mysql-8.0.3-winx64\bin;SRV*c:\symbols*http://msdl.microsoft.com/download/symbols"^
 -v -n -c "!analyze -vvvvv"
```

Note

The `^` character and newline are removed by the Windows command line processor, so be sure the spaces remain intact.

28.5.1.4 Debugging mysqld under gdb

On most systems you can also start `mysqld` from `gdb` to get more information if `mysqld` crashes.

With some older `gdb` versions on Linux you must use `run --one-thread` if you want to be able to debug `mysqld` threads. In this case, you can only have one thread active at a time. It is best to upgrade to `gdb` 5.1 because thread debugging works much better with this version!

NPTL threads (the new thread library on Linux) may cause problems while running `mysqld` under `gdb`. Some symptoms are:

- `mysqld` hangs during startup (before it writes `ready for connections`).
- `mysqld` crashes during a `pthread_mutex_lock()` or `pthread_mutex_unlock()` call.

In this case, you should set the following environment variable in the shell before starting `gdb`:

```
LD_ASSUME_KERNEL=2.4.1
export LD_ASSUME_KERNEL
```

When running `mysqld` under `gdb`, you should disable the stack trace with `--skip-stack-trace` to be able to catch segfaults within `gdb`.

Use the `--gdb` option to `mysqld` to install an interrupt handler for `SIGINT` (needed to stop `mysqld` with `^C` to set breakpoints) and disable stack tracing and core file handling.

It is very hard to debug MySQL under `gdb` if you do a lot of new connections the whole time as `gdb` does not free the memory for old threads. You can avoid this problem by starting `mysqld` with `thread_cache_size` set to a value equal to `max_connections` + 1. In most cases just using `--thread_cache_size=5'` helps a lot!

If you want to get a core dump on Linux if `mysqld` dies with a SIGSEGV signal, you can start `mysqld` with the `--core-file` option. This core file can be used to make a backtrace that may help you find out why `mysqld` died:

```
shell> gdb mysqld core
gdb>   backtrace full
gdb>   quit
```

See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.

If you are using `gdb` 4.17.x or above on Linux, you should install a `.gdb` file, with the following information, in your current directory:

```
set print sevenbit off
handle SIGUSR1 nostop noprint
handle SIGUSR2 nostop noprint
handle SIGWAITING nostop noprint
handle SIGLWP nostop noprint
handle SIGPIPE nostop
handle SIGALRM nostop
handle SIGHUP nostop
handle SIGTERM nostop noprint
```

If you have problems debugging threads with `gdb`, you should download gdb 5.x and try this instead. The new `gdb` version has very improved thread handling!

Here is an example how to debug `mysqld`:

```
shell> gdb /usr/local/libexec/mysqld
gdb> run
...
backtrace full # Do this when mysqld crashes
```

Include the preceding output in a bug report, which you can file using the instructions in Section 1.7, “How to Report Bugs or Problems”.

If `mysqld` hangs, you can try to use some system tools like `strace` or `/usr/proc/bin/pstack` to examine where `mysqld` has hung.

```
strace /tmp/log libexec/mysqld
```

If you are using the Perl `DBI` interface, you can turn on debugging information by using the `trace` method or by setting the `DBI_TRACE` environment variable.

28.5.1.5 Using a Stack Trace

On some operating systems, the error log contains a stack trace if `mysqld` dies unexpectedly. You can use this to find out where (and maybe why) `mysqld` died. See Section 5.4.2, “The Error Log”. To get a stack trace, you must not compile `mysqld` with the `-fomit-frame-pointer` option to gcc. See Section 28.5.1.1, “Compiling MySQL for Debugging”.

A stack trace in the error log looks something like this:

```
mysqld got signal 11;
Attempting backtrace. You can use the following information
to find out where mysqld died. If you see no messages after
this, something went terribly wrong...

stack_bottom = 0x41fd0110 thread_stack 0x40000
mysqld(my_print_stacktrace+0x32)[0x9da402]
mysqld(handle_segfault+0x28a)[0x6648e9]
/lib/libpthread.so.0[0x7f1a5af000f0]
/lib/libc.so.6(strcmp+0x2)[0x7f1a5a10f0f2]
mysqld(_Z21check_change_passwordP3THDPKcS2_Pcj+0x7c)[0x7412cb]
mysqld(_ZN16set_var_password5checkEP3THD+0xd0)[0x688354]
mysqld(_Z17sql_set_variablesP3THDP4ListI12set_var_baseE+0x68)[0x688494]
```

```
mysqld(_Z21mysql_execute_commandP3THD+0x41a0)[0x67a170]
mysqld(_Z11mysql_parseP3THDPKcjPS2_+0x282)[0x67f0ad]
mysqld(_Z16dispatch_command19enum_server_commandP3THDPcj+0xbb7[0x67fdf8]
mysqld(_Z10do_commandP3THD+0x24d)[0x6811b6]
mysqld(handle_one_connection+0x11c)[0x66e05e]
```

If resolution of function names for the trace fails, the trace contains less information:

```
mysqld got signal 11;
Attempting backtrace. You can use the following information
to find out where mysqld died. If you see no messages after
this, something went terribly wrong...

stack_bottom = 0x41fd0110 thread_stack 0x40000
[0x9da402]
[0x6648e9]
[0x7f1a5af000f0]
[0x7f1a5a10f0f2]
[0x7412cb]
[0x688354]
[0x688494]
[0x67a170]
[0x67f0ad]
[0x67fdf8]
[0x6811b6]
[0x66e05e]
```

In the latter case, you can use the `resolve_stack_dump` utility to determine where `mysqld` died by using the following procedure:

1. Copy the numbers from the stack trace to a file, for example `mysqld.stack`. The numbers should not include the surrounding square brackets:

   ```
   0x9da402
   0x6648e9
   0x7f1a5af000f0
   0x7f1a5a10f0f2
   0x7412cb
   0x688354
   0x688494
   0x67a170
   0x67f0ad
   0x67fdf8
   0x6811b6
   0x66e05e
   ```

2. Make a symbol file for the `mysqld` server:

   ```
   shell> nm -n libexec/mysqld > /tmp/mysqld.sym
   ```

   If `mysqld` is not linked statically, use the following command instead:

   ```
   shell> nm -D -n libexec/mysqld > /tmp/mysqld.sym
   ```

   If you want to decode C++ symbols, use the `--demangle`, if available, to `nm`. If your version of `nm` does not have this option, you will need to use the `c++filt` command after the stack dump has been produced to demangle the C++ names.

3. Execute the following command:

   ```
   shell> resolve_stack_dump -s /tmp/mysqld.sym -n mysqld.stack
   ```

   If you were not able to include demangled C++ names in your symbol file, process the `resolve_stack_dump` output using `c++filt`:

```
shell> resolve_stack_dump -s /tmp/mysqld.sym -n mysqld.stack | c++filt
```

This prints out where `mysqld` died. If that does not help you find out why `mysqld` died, you should create a bug report and include the output from the preceding command with the bug report.

However, in most cases it does not help us to have just a stack trace to find the reason for the problem. To be able to locate the bug or provide a workaround, in most cases we need to know the statement that killed `mysqld` and preferably a test case so that we can repeat the problem! See Section 1.7, “How to Report Bugs or Problems”.

Newer versions of `glibc` stack trace functions also print the address as relative to the object. On `glibc`-based systems (Linux), the trace for a crash within a plugin looks something like:

```
plugin/auth/auth_test_plugin.so(+0x9a6)[0x7ff4d11c29a6]
```

To translate the relative address (`+0x9a6`) into a file name and line number, use this command:

```
shell> addr2line -fie auth_test_plugin.so 0x9a6
auth_test_plugin
mysql-trunk/plugin/auth/test_plugin.c:65
```

The `addr2line` utility is part of the `binutils` package on Linux.

On Solaris, the procedure is similar. The Solaris `printstack()` already prints relative addresses:

```
plugin/auth/auth_test_plugin.so:0x1510
```

To translate, use this command:

```
shell> gaddr2line -fie auth_test_plugin.so 0x1510
mysql-trunk/plugin/auth/test_plugin.c:88
```

Windows already prints the address, function name and line:

```
000007FEF07E10A4 auth_test_plugin.dll!auth_test_plugin()[test_plugin.c:72]
```

28.5.1.6 Using Server Logs to Find Causes of Errors in mysqld

Note that before starting `mysqld` with the general query log enabled, you should check all your tables with `myisamchk`. See Chapter 5, *MySQL Server Administration*.

If `mysqld` dies or hangs, you should start `mysqld` with the general query log enabled. See Section 5.4.3, “The General Query Log”. When `mysqld` dies again, you can examine the end of the log file for the query that killed `mysqld`.

If you use the default general query log file, the log is stored in the database directory as *`host_name`*`.log` In most cases it is the last query in the log file that killed `mysqld`, but if possible you should verify this by restarting `mysqld` and executing the found query from the `mysql` command-line tools. If this works, you should also test all complicated queries that did not complete.

You can also try the command `EXPLAIN` on all `SELECT` statements that takes a long time to ensure that `mysqld` is using indexes properly. See Section 13.8.2, “EXPLAIN Syntax”.

You can find the queries that take a long time to execute by starting `mysqld` with the slow query log enabled. See Section 5.4.5, “The Slow Query Log”.

If you find the text `mysqld restarted` in the error log file (normally named `hostname.err`) you probably have found a query that causes `mysqld` to fail. If this happens, you should check all your

tables with `myisamchk` (see Chapter 5, *MySQL Server Administration*), and test the queries in the MySQL log files to see whether one fails. If you find such a query, try first upgrading to the newest MySQL version. If this does not help and you cannot find anything in the `mysql` mail archive, you should report the bug to a MySQL mailing list. The mailing lists are described at http://lists.mysql.com/, which also has links to online list archives.

If you have started `mysqld` with `--myisam-recover-options`, MySQL automatically checks and tries to repair `MyISAM` tables if they are marked as 'not closed properly' or 'crashed'. If this happens, MySQL writes an entry in the `hostname.err` file `'Warning: Checking table ...'` which is followed by `Warning: Repairing table` if the table needs to be repaired. If you get a lot of these errors, without `mysqld` having died unexpectedly just before, then something is wrong and needs to be investigated further. See Section 5.1.4, “Server Command Options”.

When the server detects `MyISAM` table corruption, it writes additional information to the error log, such as the name and line number of the source file, and the list of threads accessing the table. Example: `Got an error from thread_id=1, mi_dynrec.c:368`. This is useful information to include in bug reports.

It is not a good sign if `mysqld` did die unexpectedly, but in this case, you should not investigate the `Checking table...` messages, but instead try to find out why `mysqld` died.

28.5.1.7 Making a Test Case If You Experience Table Corruption

The following procedure applies to `MyISAM` tables. For information about steps to take when encountering `InnoDB` table corruption, see Section 1.7, “How to Report Bugs or Problems”.

If you encounter corrupted `MyISAM` tables or if `mysqld` always fails after some update statements, you can test whether the issue is reproducible by doing the following:

1. Stop the MySQL daemon with `mysqladmin shutdown`.
2. Make a backup of the tables to guard against the very unlikely case that the repair does something bad.
3. Check all tables with `myisamchk -s database/*.MYI`. Repair any corrupted tables with `myisamchk -r database/*table*.MYI`.
4. Make a second backup of the tables.
5. Remove (or move away) any old log files from the MySQL data directory if you need more space.
6. Start `mysqld` with the binary log enabled. If you want to find a statement that crashes `mysqld`, you should start the server with the general query log enabled as well. See Section 5.4.3, “The General Query Log”, and Section 5.4.4, “The Binary Log”.
7. When you have gotten a crashed table, stop the `mysqld` server.
8. Restore the backup.
9. Restart the `mysqld` server *without* the binary log enabled.
10. Re-execute the statements with `mysqlbinlog binary-log-file | mysql`. The binary log is saved in the MySQL database directory with the name `hostname-bin.NNNNNN`.
11. If the tables are corrupted again or you can get `mysqld` to die with the above command, you have found a reproducible bug. FTP the tables and the binary log to our bugs database using the instructions given in Section 1.7, “How to Report Bugs or Problems”. If you are a support customer, you can use the MySQL Customer Support Center (http://www.mysql.com/support/) to alert the MySQL team about the problem and have it fixed as soon as possible.

28.5.2 Debugging a MySQL Client

To be able to debug a MySQL client with the integrated debug package, you should configure MySQL with `-DWITH_DEBUG=1`. See Section 2.8.4, “MySQL Source-Configuration Options”.

Before running a client, you should set the `MYSQL_DEBUG` environment variable:

```
shell> MYSQL_DEBUG=d:t:O,/tmp/client.trace
shell> export MYSQL_DEBUG
```

This causes clients to generate a trace file in `/tmp/client.trace`.

If you have problems with your own client code, you should attempt to connect to the server and run your query using a client that is known to work. Do this by running `mysql` in debugging mode (assuming that you have compiled MySQL with debugging on):

```
shell> mysql --debug=d:t:O,/tmp/client.trace
```

This provides useful information in case you mail a bug report. See Section 1.7, “How to Report Bugs or Problems”.

If your client crashes at some 'legal' looking code, you should check that your `mysql.h` include file matches your MySQL library file. A very common mistake is to use an old `mysql.h` file from an old MySQL installation with new MySQL library.

28.5.3 The DBUG Package

The MySQL server and most MySQL clients are compiled with the DBUG package originally created by Fred Fish. When you have configured MySQL for debugging, this package makes it possible to get a trace file of what the program is doing. See Section 28.5.1.2, “Creating Trace Files”.

This section summarizes the argument values that you can specify in debug options on the command line for MySQL programs that have been built with debugging support. For more information about programming with the DBUG package, see the DBUG manual in the `dbug` directory of MySQL source distributions. It's best to use a recent distribution to get the most updated DBUG manual.

The DBUG package can be used by invoking a program with the `--debug[=debug_options]` or `-#[debug_options]` option. If you specify the `--debug` or `-#` option without a *`debug_options`* value, most MySQL programs use a default value. The server default is `d:t:i:o,/tmp/mysqld.trace` on Unix and `d:t:i:O,\mysqld.trace` on Windows. The effect of this default is:

- `d`: Enable output for all debug macros
- `t`: Trace function calls and exits
- `i`: Add PID to output lines
- `o,/tmp/mysqld.trace`, `O,\mysqld.trace`: Set the debug output file.

Most client programs use a default *`debug_options`* value of `d:t:o,/tmp/`*`program_name`*`.trace`, regardless of platform.

Here are some example debug control strings as they might be specified on a shell command line:

```
--debug=d:t
--debug=d:f,main,subr1:F:L:t,20
--debug=d,input,output,files:n
--debug=d:t:i:O,\\mysqld.trace
```

For `mysqld`, it is also possible to change DBUG settings at runtime by setting the `debug` system variable. This variable has global and session values:

```
mysql> SET GLOBAL debug = 'debug_options';
mysql> SET SESSION debug = 'debug_options';
```

Changes at runtime require the SYSTEM_VARIABLES_ADMIN or SUPER privilege, even for the session value.

The *debug_options* value is a sequence of colon-separated fields:

```
field_1:field_2:...:field_N
```

Each field within the value consists of a mandatory flag character, optionally preceded by a + or - character, and optionally followed by a comma-delimited list of modifiers:

```
[+|-]flag[,modifier,modifier,...,modifier]
```

The following table describes the permitted flag characters. Unrecognized flag characters are silently ignored.

Flag	Description
d	Enable output from DBUG_*XXX* macros for the current state. May be followed by a list of keywords, which enables output only for the DBUG macros with that keyword. An empty list of keywords enables output for all macros. In MySQL, common debug macro keywords to enable are enter, exit, error, warning, info, and loop.
D	Delay after each debugger output line. The argument is the delay, in tenths of seconds, subject to machine capabilities. For example, D,20 specifies a delay of two seconds.
f	Limit debugging, tracing, and profiling to the list of named functions. An empty list enables all functions. The appropriate d or t flags must still be given; this flag only limits their actions if they are enabled.
F	Identify the source file name for each line of debug or trace output.
i	Identify the process with the PID or thread ID for each line of debug or trace output.
L	Identify the source file line number for each line of debug or trace output.
n	Print the current function nesting depth for each line of debug or trace output.
N	Number each line of debug output.
o	Redirect the debugger output stream to the specified file. The default output is stderr.
O	Like o, but the file is really flushed between each write. When needed, the file is closed and reopened between each write.
p	Limit debugger actions to specified processes. A process must be identified with the DBUG_PROCESS macro and match one in the list for debugger actions to occur.
P	Print the current process name for each line of debug or trace output.
r	When pushing a new state, do not inherit the previous state's function nesting level. Useful when the output is to start at the left margin.
S	Do function _sanity(_file_,_line_) at each debugged function until _sanity() returns something that differs from 0.
t	Enable function call/exit trace lines. May be followed by a list (containing only one modifier) giving a numeric maximum trace level, beyond which no output occurs for either debugging or tracing macros. The default is a compile time option.

The leading + or - character and trailing list of modifiers are used for flag characters such as d or f that can enable a debug operation for all applicable modifiers or just some of them:

- With no leading + or -, the flag value is set to exactly the modifier list as given.

- With a leading + or -, the modifiers in the list are added to or subtracted from the current modifier list.

The following examples show how this works for the d flag. An empty d list enabled output for all debug macros. A nonempty list enables output only for the macro keywords in the list.

These statements set the d value to the modifier list as given:

```
mysql> SET debug = 'd';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| d       |
+---------+
mysql> SET debug = 'd,error,warning';
mysql> SELECT @@debug;
+-----------------+
| @@debug         |
+-----------------+
| d,error,warning |
+-----------------+
```

A leading + or - adds to or subtracts from the current d value:

```
mysql> SET debug = '+d,loop';
mysql> SELECT @@debug;
+----------------------+
| @@debug              |
+----------------------+
| d,error,warning,loop |
+----------------------+
mysql> SET debug = '-d,error,loop';
mysql> SELECT @@debug;
+-----------+
| @@debug   |
+-----------+
| d,warning |
+-----------+
```

Adding to “all macros enabled” results in no change:

```
mysql> SET debug = 'd';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| d       |
+---------+
mysql> SET debug = '+d,loop';
mysql> SELECT @@debug;
+---------+
| @@debug |
+---------+
| d       |
+---------+
```

Disabling all enabled macros disables the d flag entirely:

```
mysql> SET debug = 'd,error,loop';
mysql> SELECT @@debug;
+--------------+
| @@debug      |
+--------------+
| d,error,loop |
+--------------+
mysql> SET debug = '-d,error,loop';
mysql> SELECT @@debug;
```

```
+---------+
| @@debug |
+---------+
|         |
+---------+
```

Chapter 29 MySQL Enterprise Edition

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MySQL Enterprise Edition is a commercial product. Like MySQL Community Edition, MySQL Enterprise Edition includes MySQL Server, a fully integrated transaction-safe, ACID-compliant database with full commit, rollback, crash-recovery, and row-level locking capabilities. In addition, MySQL Enterprise Edition includes the following components designed to provide monitoring and online backup, as well as improved security and scalability:

The following sections briefly discuss each of these components and indicate where to find more detailed information. To learn more about commercial products, see http://www.mysql.com/products/.

- MySQL Enterprise Monitor
- MySQL Enterprise Backup
- MySQL Enterprise Security
- MySQL Enterprise Encryption
- MySQL Enterprise Audit
- MySQL Enterprise Firewall
- MySQL Enterprise Thread Pool

29.1 MySQL Enterprise Monitor Overview

MySQL Enterprise Monitor is an enterprise monitoring system for MySQL that keeps an eye on your MySQL servers, notifies you of potential issues and problems, and advises you how to fix the issues. MySQL Enterprise Monitor can monitor all kinds of configurations, from a single MySQL server that is important to your business, all the way up to a huge farm of MySQL servers powering a busy web site.

The following discussion briefly summarizes the basic components that make up the MySQL Enterprise Monitor product. For more information, see the MySQL Enterprise Monitor manual, available at http://dev.mysql.com/doc/mysql-monitor/en/.

MySQL Enterprise Monitor components can be installed in various configurations depending on your database and network topology, to give you the best combination of reliable and responsive monitoring data, with minimal overhead on the database server machines. A typical MySQL Enterprise Monitor installation consists of:

- One or more MySQL servers to monitor. MySQL Enterprise Monitor can monitor both Community and Enterprise MySQL server releases.
- A MySQL Enterprise Monitor Agent for each monitored host.
- A single MySQL Enterprise Service Manager, which collates information from the agents and provides the user interface to the collected data.

MySQL Enterprise Monitor is designed to monitor one or more MySQL servers. The monitoring information is collected by using an agent, *MySQL Enterprise Monitor Agent*. The agent communicates with the hosts and MySQL servers that it monitors, collecting variables, status and health information, and sending this information to the MySQL Enterprise Service Manager.

The information collected by the agent about each MySQL server and host you are monitoring is sent to the *MySQL Enterprise Service Manager*. This server collates all of the information from the agents. As it collates the information sent by the agents, the MySQL Enterprise Service Manager continually tests the collected data, comparing the status of the server to reasonable values. When thresholds are reached, the server can trigger an event (including an alarm and notification) to highlight a potential issue, such as low memory, high CPU usage, or more complex conditions such insufficient buffer sizes and status information. We call each test, with its associated threshold value, a *rule*.

These rules, and the alarms and notifications, are each known as a *MySQL Enterprise Advisors*. Advisors form a critical part of the MySQL Enterprise Service Manager, as they provide warning information and troubleshooting advice about potential problems.

The MySQL Enterprise Service Manager includes a web server, and you interact with it through any web browser. This interface, the MySQL Enterprise Monitor User Interface, displays all of the information collected by the agents, and lets you view all of your servers and their current status as a group or individually. You control and configure all aspects of the service using the MySQL Enterprise Monitor User Interface.

The information supplied by the MySQL Enterprise Monitor Agent processes also includes statistical and query information, which you can view in the form of graphs. For example, you can view aspects such as server load, query numbers, or index usage information as a graph over time. The graph lets you pinpoint problems or potential issues on your server, and can help diagnose the impact from database or external problems (such as external system or network failure) by examining the data from a specific time interval.

The MySQL Enterprise Monitor Agent can also be configured to collect detailed information about the queries executed on your server, including the row counts and performance times for executing each query. You can correlate the detailed query data with the graphical information to identify which queries were executing when you experienced a particularly high load, index or other issue. The query data is supported by a system called Query Analyzer, and the data can be presented in different ways depending on your needs.

29.2 MySQL Enterprise Backup Overview

MySQL Enterprise Backup performs hot backup operations for MySQL databases. The product is architected for efficient and reliable backups of tables created by the InnoDB storage engine. For completeness, it can also back up tables from MyISAM and other storage engines.

The following discussion briefly summarizes MySQL Enterprise Backup. For more information, see the MySQL Enterprise Backup manual, available at http://dev.mysql.com/doc/mysql-enterprise-backup/en/.

Hot backups are performed while the database is running and applications are reading and writing to it. This type of backup does not block normal database operations, and it captures even changes that occur while the backup is happening. For these reasons, hot backups are desirable when your database “grows up” -- when the data is large enough that the backup takes significant time, and when your data is important enough to your business that you must capture every last change, without taking your application, web site, or web service offline.

MySQL Enterprise Backup does a hot backup of all tables that use the InnoDB storage engine. For tables using MyISAM or other non-InnoDB storage engines, it does a “warm” backup, where the database continues to run, but those tables cannot be modified while being backed up. For efficient backup operations, you can designate InnoDB as the default storage engine for new tables, or convert existing tables to use the InnoDB storage engine.

29.3 MySQL Enterprise Security Overview

MySQL Enterprise Edition provides plugins that implement authentication using external services:

- MySQL Enterprise Edition includes an authentication plugin that enables MySQL Server to use PAM (Pluggable Authentication Modules) to authenticate MySQL users. PAM enables a system to use a standard interface to access various kinds of authentication methods, such as Unix passwords or an LDAP directory. For more information, see PAM Pluggable Authentication.
- MySQL Enterprise Edition includes an authentication plugin that performs external authentication on Windows, enabling MySQL Server to use native Windows services to authenticate client connections. Users who have logged in to Windows can connect from MySQL client programs to the server based on the information in their environment without specifying an additional password. For more information, see Windows Pluggable Authentication.
- MySQL Enterprise Edition includes a keyring plugin that uses Oracle Key Vault as a back end for keyring storage. For more information, see Section 6.5.4, “The MySQL Keyring”.

For other related Enterprise security features, see Section 29.4, “MySQL Enterprise Encryption Overview”.

29.4 MySQL Enterprise Encryption Overview

MySQL Enterprise Edition includes a set of encryption functions based on the OpenSSL library that expose OpenSSL capabilities at the SQL level. These functions enable Enterprise applications to perform the following operations:

- Implement added data protection using public-key asymmetric cryptography
- Create public and private keys and digital signatures
- Perform asymmetric encryption and decryption
- Use cryptographic hashing for digital signing and data verification and validation

For more information, see MySQL Enterprise Encryption Functions.

For other related Enterprise security features, see Section 29.3, “MySQL Enterprise Security Overview”.

29.5 MySQL Enterprise Audit Overview

MySQL Enterprise Edition includes MySQL Enterprise Audit, implemented using a server plugin. MySQL Enterprise Audit uses the open MySQL Audit API to enable standard, policy-based monitoring and logging of connection and query activity executed on specific MySQL servers. Designed to meet the Oracle audit specification, MySQL Enterprise Audit provides an out of box, easy to use auditing and compliance solution for applications that are governed by both internal and external regulatory guidelines.

When installed, the audit plugin enables MySQL Server to produce a log file containing an audit record of server activity. The log contents include when clients connect and disconnect, and what actions they perform while connected, such as which databases and tables they access.

For more information, see MySQL Enterprise Audit.

29.6 MySQL Enterprise Firewall Overview

MySQL Enterprise Edition includes MySQL Enterprise Firewall, an application-level firewall that enables database administrators to permit or deny SQL statement execution based on matching against whitelists of accepted statement patterns. This helps harden MySQL Server against attacks such as SQL injection or attempts to exploit applications by using them outside of their legitimate query workload characteristics.

Each MySQL account registered with the firewall has its own statement whitelist, enabling protection to be tailored per account. For a given account, the firewall can operate in recording or protecting mode, for training in the accepted statement patterns or protection against unacceptable statements.

For more information, see MySQL Enterprise Firewall.

29.7 MySQL Enterprise Thread Pool Overview

MySQL Enterprise Edition includes MySQL Enterprise Thread Pool, implemented using a server plugin. The default thread-handling model in MySQL Server executes statements using one thread per client connection. As more clients connect to the server and execute statements, overall performance degrades. In MySQL Enterprise Edition, a thread pool plugin provides an alternative thread-handling model designed to reduce overhead and improve performance. The plugin implements a thread pool that increases server performance by efficiently managing statement execution threads for large numbers of client connections.

For more information, see MySQL Enterprise Thread Pool.

Chapter 30 MySQL Workbench

MySQL Workbench provides a graphical tool for working with MySQL servers and databases. MySQL Workbench fully supports MySQL versions 5.1 and above.

The following discussion briefly describes MySQL Workbench capabilities. For more information, see the MySQL Workbench manual, available at http://dev.mysql.com/doc/workbench/en/.

MySQL Workbench provides five main areas of functionality:

- **SQL Development**: Enables you to create and manage connections to database servers. As well as enabling you to configure connection parameters, MySQL Workbench provides the capability to execute SQL queries on the database connections using the built-in SQL Editor. This functionality replaces that previously provided by the Query Browser standalone application.
- **Data Modeling**: Enables you to create models of your database schema graphically, reverse and forward engineer between a schema and a live database, and edit all aspects of your database using the comprehensive Table Editor. The Table Editor provides easy-to-use facilities for editing Tables, Columns, Indexes, Triggers, Partitioning, Options, Inserts and Privileges, Routines and Views.
- **Server Administration**: Enables you to create and administer server instances.
- **Data Migration**: Allows you to migrate from Microsoft SQL Server, Sybase ASE, SQLite, SQL Anywhere, PostreSQL, and other RDBMS tables, objects and data to MySQL. Migration also supports migrating from earlier versions of MySQL to the latest releases.
- **MySQL Enterprise Support**: Support for Enterprise products such as MySQL Enterprise Backup and MySQL Audit.

MySQL Workbench is available in two editions, the Community Edition and the Commercial Edition. The Community Edition is available free of charge. The Commercial Edition provides additional Enterprise features, such as database documentation generation, at low cost.

Appendix A MySQL 8.0 Frequently Asked Questions

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A.1 MySQL 8.0 FAQ: General

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A.1.1. Which version of MySQL is production-ready (GA)?

MySQL 5.7 and MySQL 5.6 are supported for production use.

MySQL 5.7 achieved General Availability (GA) status with MySQL 5.7.9, which was released for production use on 21 October 2015.

MySQL 5.6 achieved General Availability (GA) status with MySQL 5.6.10, which was released for production use on 5 February 2013.

MySQL 5.5 achieved General Availability (GA) status with MySQL 5.5.8, which was released for production use on 3 December 2010. The MySQL 5.5 series is no longer current, but still supported in production.

MySQL 5.1 achieved General Availability (GA) status with MySQL 5.1.30, which was released for production use on 14 November 2008. Active development for MySQL 5.1 has ended.

MySQL 5.0 achieved General Availability (GA) status with MySQL 5.0.15, which was released for production use on 19 October 2005. Active development for MySQL 5.0 has ended.

A.1.2. What is the state of development (non-GA) versions?

MySQL follows a milestone release model that introduces pre-production-quality features and stabilizes them to release quality (see http://dev.mysql.com/doc/mysql-development-cycle/en/index.html). This process then repeats, so releases cycle between pre-production and release quality status. Please check the change logs to identify the status of a given release.

MySQL 8.0, a successor to MySQL 5.7, is being actively developed using the milestone release methodology described above.

A.1.3. Why did MySQL version numbering skip versions 6 and 7 and go straight to 8.0?

Due to the many new and important features we were introducing in this MySQL version, we decided to start a fresh new series. As the series numbers 6 and 7 had actually been used before by MySQL, we went to 8.0.

A.1.4. Can MySQL 8.0 do subqueries?

Yes. See Section 13.2.11, “Subquery Syntax”.

A.1.5. Can MySQL 8.0 perform multiple-table inserts, updates, and deletes?

Yes. For the syntax required to perform multiple-table updates, see Section 13.2.12, “UPDATE Syntax”; for that required to perform multiple-table deletes, see Section 13.2.2, “DELETE Syntax”.

A multiple-table insert can be accomplished using a trigger whose `FOR EACH ROW` clause contains multiple `INSERT` statements within a `BEGIN ... END` block. See Section 23.3, “Using Triggers”.

A.1.6. Does MySQL 8.0 have a Query Cache? Does it work on Server, Instance or Database?

Yes. The query cache operates on the server level, caching complete result sets matched with the original query string. If an exactly identical query is made (which often happens, particularly in web applications), no parsing or execution is necessary; the result is sent directly from the cache. Various tuning options are available. See Section 8.10.3, “The MySQL Query Cache”.

A.1.7. Does MySQL 8.0 have Sequences?

No. However, MySQL has an `AUTO_INCREMENT` system, which in MySQL 8.0 can also handle inserts in a multi-master replication setup. With the `auto_increment_increment` and `auto_increment_offset` system variables, you can set each server to generate auto-increment values that don't conflict with other servers. The `auto_increment_increment` value should be greater than the number of servers, and each server should have a unique offset.

A.1.8. Does MySQL 8.0 have a `NOW()` function with fractions of seconds?

Yes, see Section 11.3.6, “Fractional Seconds in Time Values”.

A.1.9. Does MySQL 8.0 work with multi-core processors?

Yes. MySQL is fully multi-threaded, and will make use of multiple CPUs, provided that the operating system supports them.

A.1.10 Why do I see multiple processes for `mysqld`?

When using LinuxThreads, you should see a minimum of three `mysqld` processes running. These are in fact threads. There is one thread for the LinuxThreads manager, one thread to handle connections, and one thread to handle alarms and signals.

A.1.11 Can MySQL 8.0 perform ACID transactions?

Yes. All current MySQL versions support transactions. The `InnoDB` storage engine offers full ACID transactions with row-level locking, multi-versioning, nonlocking repeatable reads, and all four SQL standard isolation levels.

The NDB storage engine supports the READ COMMITTED transaction isolation level only.

A.2 MySQL 8.0 FAQ: Storage Engines

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A.2.1. Where can I obtain complete documentation for MySQL storage engines?

See Chapter 16, *Alternative Storage Engines*. That chapter contains information about all MySQL storage engines except for the InnoDB storage engine and the NDB storage engine (used for MySQL Cluster). InnoDB is covered in Chapter 15, *The InnoDB Storage Engine*. NDB is covered in MySQL NDB Cluster 7.5 and NDB Cluster 7.6.

A.2.2. Are there any new storage engines in MySQL 8.0?

No. InnoDB is the default storage engine for new tables. See Section 15.1, “Introduction to InnoDB” for details.

A.2.3. Have any storage engines been removed in MySQL 8.0?

The PARTITION storage engine plugin which provided partitioning support is replaced by a native partitioning handler. As part of this change, the server can no longer be built using -DWITH_PARTITION_STORAGE_ENGINE. partition is also no longer displayed in the output of SHOW PLUGINS, or shown in the INFORMATION_SCHEMA.PLUGINS table.

In order to support partitioning of a given table, the storage engine used for the table must now provide its own (“native”) partitioning handler. InnoDB is the only storage engine supported in MySQL 8.0 that includes a native partitioning handler. An attempt to create partitioned tables in MySQL 8.0 using any other storage engine fails. (The NDB storage engine used by MySQL Cluster also provides its own partitioning handler, but is currently not supported by MySQL 8.0.)

A.2.4. Can I prevent the use of a particular storage engine?

Yes. The disabled_storage_engines configuration option defines which storage engines cannot be used to create tables or tablespaces. By default, disabled_storage_engines is empty (no engines disabled), but it can be set to a comma-separated list of one or more engines.

A.2.5. Is there an advantage to using the InnoDB storage engine exclusively, as opposed to a combination of InnoDB and non-InnoDB storage engines?

Yes. Using InnoDB tables exclusively can simplify backup and recovery operations. MySQL Enterprise Backup does a hot backup of all tables that use the InnoDB storage engine. For tables using MyISAM or other non-InnoDB storage engines, it does a “warm” backup, where the database continues to run, but those tables cannot be modified while being backed up. See Section 29.2, “MySQL Enterprise Backup Overview”.

A.2.6. What are the unique benefits of the ARCHIVE storage engine?

The ARCHIVE storage engine stores large amounts of data without indexes; it has a small footprint, and performs selects using table scans. See Section 16.5, “The ARCHIVE Storage Engine”, for details.

A.3 MySQL 8.0 FAQ: Server SQL Mode

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A.3.1. What are server SQL modes?

Server SQL modes define what SQL syntax MySQL should support and what kind of data validation checks it should perform. This makes it easier to use MySQL in different environments and to use MySQL together with other database servers. The MySQL Server apply these modes individually to different clients. For more information, see Section 5.1.8, “Server SQL Modes”.

A.3.2. How many server SQL modes are there?

Each mode can be independently switched on and off. See Section 5.1.8, “Server SQL Modes”, for a complete list of available modes.

A.3.3. How do you determine the server SQL mode?

You can set the default SQL mode (for `mysqld` startup) with the `--sql-mode` option. Using the statement `SET [GLOBAL|SESSION] sql_mode='modes'`, you can change the settings from within a connection, either locally to the connection, or to take effect globally. You can retrieve the current mode by issuing a `SELECT @@sql_mode` statement.

A.3.4. Is the mode dependent on the database or connection?

A mode is not linked to a particular database. Modes can be set locally to the session (connection), or globally for the server. you can change these settings using `SET [GLOBAL|SESSION] sql_mode='modes'`.

A.3.5. Can the rules for strict mode be extended?

When we refer to *strict mode*, we mean a mode where at least one of the modes `TRADITIONAL`, `STRICT_TRANS_TABLES`, or `STRICT_ALL_TABLES` is enabled. Options can be combined, so you can add restrictions to a mode. See Section 5.1.8, “Server SQL Modes”, for more information.

A.3.6. Does strict mode impact performance?

The intensive validation of input data that some settings requires more time than if the validation is not done. While the performance impact is not that great, if you do not require such validation (perhaps your application already handles all of this), then MySQL gives you the option of leaving strict mode disabled. However, if you do require it, strict mode can provide such validation.

A.3.7. What is the default server SQL mode when MySQL 8.0 is installed?

The default SQL mode in MySQL 8.0 includes these modes: `ONLY_FULL_GROUP_BY`, `STRICT_TRANS_TABLES`, `NO_ZERO_IN_DATE`, `NO_ZERO_DATE`, `ERROR_FOR_DIVISION_BY_ZERO`, `NO_AUTO_CREATE_USER`, and `NO_ENGINE_SUBSTITUTION`.

A.4 MySQL 8.0 FAQ: Stored Procedures and Functions

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A.4.1. Does MySQL 8.0 support stored procedures and functions?

Yes. MySQL 8.0 supports two types of stored routines, stored procedures and stored functions.

A.4.2. Where can I find documentation for MySQL stored procedures and stored functions?

See Section 23.2, “Using Stored Routines (Procedures and Functions)”.

A.4.3. Is there a discussion forum for MySQL stored procedures?

Yes. See http://forums.mysql.com/list.php?98.

A.4.4. Where can I find the ANSI SQL 2003 specification for stored procedures?

Unfortunately, the official specifications are not freely available (ANSI makes them available for purchase). However, there are books, such as *SQL-99 Complete, Really* by Peter Gulutzan and Trudy Pelzer, that provide a comprehensive overview of the standard, including coverage of stored procedures.

A.4.5. How do you manage stored routines?

It is always good practice to use a clear naming scheme for your stored routines. You can manage stored procedures with `CREATE [FUNCTION|PROCEDURE]`, `ALTER [FUNCTION|PROCEDURE]`, `DROP [FUNCTION|PROCEDURE]`, and `SHOW CREATE [FUNCTION|PROCEDURE]`. You can obtain information about existing stored procedures using the `ROUTINES` table in the `INFORMATION_SCHEMA` database (see Section 24.18, “The INFORMATION_SCHEMA ROUTINES Table”).

A.4.6. Is there a way to view all stored procedures and stored functions in a given database?

Yes. For a database named *`dbname`*, use this query on the `INFORMATION_SCHEMA.ROUTINES` table:

```
SELECT ROUTINE_TYPE, ROUTINE_NAME
    FROM INFORMATION_SCHEMA.ROUTINES
    WHERE ROUTINE_SCHEMA='dbname';
```

For more information, see Section 24.18, “The INFORMATION_SCHEMA ROUTINES Table”.

The body of a stored routine can be viewed using `SHOW CREATE FUNCTION` (for a stored function) or `SHOW CREATE PROCEDURE` (for a stored procedure). See Section 13.7.5.9, “SHOW CREATE PROCEDURE Syntax”, for more information.

A.4.7. Where are stored procedures stored?

Stored procedures are stored in the `mysql.routines` and `mysql.parameters` tables, which are part of the data dictionary. You cannot access these tables directly. Instead, query the `INFORMATION_SCHEMA ROUTINES` and `PARAMETERS` tables. See Section 24.18, “The INFORMATION_SCHEMA ROUTINES Table”, and Section 24.12, “The INFORMATION_SCHEMA PARAMETERS Table”.

You can also use `SHOW CREATE FUNCTION` to obtain information about stored functions, and `SHOW CREATE PROCEDURE` to obtain information about stored procedures. See Section 13.7.5.9, “SHOW CREATE PROCEDURE Syntax”.

A.4.8. Is it possible to group stored procedures or stored functions into packages?

No. This is not supported in MySQL 8.0.

A.4.9. Can a stored procedure call another stored procedure?

Yes.

A.4.10 Can a stored procedure call a trigger?

A stored procedure can execute an SQL statement, such as an `UPDATE`, that causes a trigger to activate.

A.4.11 Can a stored procedure access tables?

Yes. A stored procedure can access one or more tables as required.

A.4.12 Do stored procedures have a statement for raising application errors?

Yes. MySQL 8.0 implements the SQL standard `SIGNAL` and `RESIGNAL` statements. See Section 13.6.7, “Condition Handling”.

A.4.13 Do stored procedures provide exception handling?

MySQL implements `HANDLER` definitions according to the SQL standard. See Section 13.6.7.2, “DECLARE ... HANDLER Syntax”, for details.

A.4.14 Can MySQL 8.0 stored routines return result sets?

Stored procedures can, but stored functions cannot. If you perform an ordinary `SELECT` inside a stored procedure, the result set is returned directly to the client. You need to use the MySQL 4.1 (or higher) client/server protocol for this to work. This means that, for example, in PHP, you need to use the `mysqli` extension rather than the old `mysql` extension.

A.4.15 Is `WITH RECOMPILE` supported for stored procedures?

Not in MySQL 8.0.

A.4.16 Is there a MySQL equivalent to using `mod_plsql` as a gateway on Apache to talk directly to a stored procedure in the database?

There is no equivalent in MySQL 8.0.

A.4.17 Can I pass an array as input to a stored procedure?

Not in MySQL 8.0.

A.4.18 Can I pass a cursor as an `IN` parameter to a stored procedure?

In MySQL 8.0, cursors are available inside stored procedures only.

A.4.19 Can I return a cursor as an `OUT` parameter from a stored procedure?

In MySQL 8.0, cursors are available inside stored procedures only. However, if you do not open a cursor on a `SELECT`, the result will be sent directly to the client. You can also `SELECT INTO` variables. See Section 13.2.10, “SELECT Syntax”.

A.4.20 Can I print out a variable's value within a stored routine for debugging purposes?

Yes, you can do this in a *stored procedure*, but not in a stored function. If you perform an ordinary `SELECT` inside a stored procedure, the result set is returned directly to the client. You will need to use the MySQL 4.1 (or above) client/server protocol for this to work. This means that, for example, in PHP, you need to use the `mysqli` extension rather than the old `mysql` extension.

A.4.21 Can I commit or roll back transactions inside a stored procedure?

Yes. However, you cannot perform transactional operations within a stored function.

A.4.22 Do MySQL 8.0 stored procedures and functions work with replication?

Yes, standard actions carried out in stored procedures and functions are replicated from a master MySQL server to a slave server. There are a few limitations that are described in detail in Section 23.7, “Binary Logging of Stored Programs”.

A.4.23 Are stored procedures and functions created on a master server replicated to a slave?

Yes, creation of stored procedures and functions carried out through normal DDL statements on a master server are replicated to a slave, so the objects will exist on both servers. `ALTER` and `DROP` statements for stored procedures and functions are also replicated.

A.4.24 How are actions that take place inside stored procedures and functions replicated?

MySQL records each DML event that occurs in a stored procedure and replicates those individual actions to a slave server. The actual calls made to execute stored procedures are not replicated.

Stored functions that change data are logged as function invocations, not as the DML events that occur inside each function.

A.4.25 Are there special security requirements for using stored procedures and functions together with replication?

Yes. Because a slave server has authority to execute any statement read from a master's binary log, special security constraints exist for using stored functions with replication. If replication or binary logging in general (for the purpose of point-in-time recovery) is active, then MySQL DBAs have two security options open to them:

1. Any user wishing to create stored functions must be granted the `SUPER` privilege.
2. Alternatively, a DBA can set the `log_bin_trust_function_creators` system variable to 1, which enables anyone with the standard `CREATE ROUTINE` privilege to create stored functions.

A.4.26 What limitations exist for replicating stored procedure and function actions?

Nondeterministic (random) or time-based actions embedded in stored procedures may not replicate properly. By their very nature, randomly produced results are not predictable and cannot be exactly reproduced, and therefore, random actions replicated to a slave will not mirror those performed on a master. Declaring stored functions to be `DETERMINISTIC` or setting the `log_bin_trust_function_creators` system variable to 0 will not allow random-valued operations to be invoked.

In addition, time-based actions cannot be reproduced on a slave because the timing of such actions in a stored procedure is not reproducible through the binary log used for replication. It records only DML events and does not factor in timing constraints.

Finally, nontransactional tables for which errors occur during large DML actions (such as bulk inserts) may experience replication issues in that a master may be partially updated from DML activity, but no updates are done to the slave because of the errors that occurred. A workaround is for a function's DML actions to be carried out with the `IGNORE` keyword so that updates on the master that cause errors are ignored and updates that do not cause errors are replicated to the slave.

A.4.27 Do the preceding limitations affect the ability of MySQL to do point-in-time recovery?

The same limitations that affect replication do affect point-in-time recovery.

A.4.28 What is being done to correct the aforementioned limitations?

You can choose either statement-based replication or row-based replication. The original replication implementation is based on statement-based binary logging. Row-based binary logging resolves the limitations mentioned earlier.

Mixed replication is also available (by starting the server with `--binlog-format=mixed`). This hybrid, “smart” form of replication “knows” whether statement-level replication can safely be used, or row-level replication is required.

For additional information, see Section 18.2.1, “Replication Formats”.

A.5 MySQL 8.0 FAQ: Triggers

A.5.1 Where can I find the documentation for MySQL 8.0 triggers? 3396
A.5.2 Is there a discussion forum for MySQL Triggers? 3396
A.5.3 Does MySQL 8.0 have statement-level or row-level triggers? 3396
A.5.4 Are there any default triggers? 3396
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A.5.6 Is there a way to view all triggers in a given database? 3397
A.5.7 Where are triggers stored? 3397
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A.5.10 Can a table have multiple triggers with the same trigger event and action time? 3397
A.5.11 Can triggers call an external application through a UDF? 3397
A.5.12 Is it possible for a trigger to update tables on a remote server? 3397
A.5.13 Do triggers work with replication? 3397
A.5.14 How are actions carried out through triggers on a master replicated to a slave? 3398

A.5.1. Where can I find the documentation for MySQL 8.0 triggers?

See Section 23.3, “Using Triggers”.

A.5.2. Is there a discussion forum for MySQL Triggers?

Yes. It is available at http://forums.mysql.com/list.php?99.

A.5.3. Does MySQL 8.0 have statement-level or row-level triggers?

In MySQL 8.0, all triggers are `FOR EACH ROW`; that is, the trigger is activated for each row that is inserted, updated, or deleted. MySQL 8.0 does not support triggers using `FOR EACH STATEMENT`.

A.5.4. Are there any default triggers?

Not explicitly. MySQL does have specific special behavior for some `TIMESTAMP` columns, as well as for columns which are defined using `AUTO_INCREMENT`.

A.5.5. How are triggers managed in MySQL?

In MySQL 8.0, triggers can be created using the `CREATE TRIGGER` statement, and dropped using `DROP TRIGGER`. See Section 13.1.17, “CREATE TRIGGER Syntax”, and Section 13.1.27, “DROP TRIGGER Syntax”, for more about these statements.

Information about triggers can be obtained by querying the `INFORMATION_SCHEMA.TRIGGERS` table. See Section 24.28, “The INFORMATION_SCHEMA TRIGGERS Table”.

A.5.6. Is there a way to view all triggers in a given database?

Yes. You can obtain a listing of all triggers defined on database `dbname` using a query on the `INFORMATION_SCHEMA.TRIGGERS` table such as the one shown here:

```
SELECT TRIGGER_NAME, EVENT_MANIPULATION, EVENT_OBJECT_TABLE, ACTION_STATEMENT
    FROM INFORMATION_SCHEMA.TRIGGERS
    WHERE TRIGGER_SCHEMA='dbname';
```

For more information about this table, see Section 24.28, “The INFORMATION_SCHEMA TRIGGERS Table”.

You can also use the `SHOW TRIGGERS` statement, which is specific to MySQL. See Section 13.7.5.38, “SHOW TRIGGERS Syntax”.

A.5.7. Where are triggers stored?

Triggers are stored in the `mysql.triggers` system table, which is part of the data dictionary.

A.5.8. Can a trigger call a stored procedure?

Yes.

A.5.9. Can triggers access tables?

A trigger can access both old and new data in its own table. A trigger can also affect other tables, but it is not permitted to modify a table that is already being used (for reading or writing) by the statement that invoked the function or trigger.

A.5.10. Can a table have multiple triggers with the same trigger event and action time?

In MySQL 8.0, it is possible to define multiple triggers for a given table that have the same trigger event and action time. For example, you can have two `BEFORE UPDATE` triggers for a table. By default, triggers that have the same trigger event and action time activate in the order they were created. To affect trigger order, specify a clause after `FOR EACH ROW` that indicates `FOLLOWS` or `PRECEDES` and the name of an existing trigger that also has the same trigger event and action time. With `FOLLOWS`, the new trigger activates after the existing trigger. With `PRECEDES`, the new trigger activates before the existing trigger.

A.5.11. Can triggers call an external application through a UDF?

Yes. For example, a trigger could invoke the `sys_exec()` UDF.

A.5.12. Is it possible for a trigger to update tables on a remote server?

Yes. A table on a remote server could be updated using the `FEDERATED` storage engine. (See Section 16.8, “The FEDERATED Storage Engine”).

A.5.13. Do triggers work with replication?

Yes. However, the way in which they work depends whether you are using MySQL's “classic” statement-based replication available in all versions of MySQL, or the row-based replication format introduced in MySQL 5.1.

When using statement-based replication, triggers on the slave are executed by statements that are executed on the master (and replicated to the slave).

When using row-based replication, triggers are not executed on the slave due to statements that were run on the master and then replicated to the slave. Instead, when using row-based replication, the changes caused by executing the trigger on the master are applied on the slave.

For more information, see Section 18.4.1.35, “Replication and Triggers”.

A.5.14 How are actions carried out through triggers on a master replicated to a slave?

Again, this depends on whether you are using statement-based or row-based replication.

Statement-based replication. First, the triggers that exist on a master must be re-created on the slave server. Once this is done, the replication flow works as any other standard DML statement that participates in replication. For example, consider a table `EMP` that has an `AFTER` insert trigger, which exists on a master MySQL server. The same `EMP` table and `AFTER` insert trigger exist on the slave server as well. The replication flow would be:

1. An `INSERT` statement is made to `EMP`.
2. The `AFTER` trigger on `EMP` activates.
3. The `INSERT` statement is written to the binary log.
4. The replication slave picks up the `INSERT` statement to `EMP` and executes it.
5. The `AFTER` trigger on `EMP` that exists on the slave activates.

Row-based replication. When you use row-based replication, the changes caused by executing the trigger on the master are applied on the slave. However, the triggers themselves are not actually executed on the slave under row-based replication. This is because, if both the master and the slave applied the changes from the master and, in addition, the trigger causing these changes were applied on the slave, the changes would in effect be applied twice on the slave, leading to different data on the master and the slave.

In most cases, the outcome is the same for both row-based and statement-based replication. However, if you use different triggers on the master and slave, you cannot use row-based replication. (This is because the row-based format replicates the changes made by triggers executing on the master to the slaves, rather than the statements that caused the triggers to execute, and the corresponding triggers on the slave are not executed.) Instead, any statements causing such triggers to be executed must be replicated using statement-based replication.

For more information, see Section 18.4.1.35, “Replication and Triggers”.

A.6 MySQL 8.0 FAQ: Views

A.6.1 Where can I find documentation covering MySQL Views? 3398
A.6.2 Is there a discussion forum for MySQL Views? 3398
A.6.3 What happens to a view if an underlying table is dropped or renamed? 3399
A.6.4 Does MySQL 8.0 have table snapshots? 3399
A.6.5 Does MySQL 8.0 have materialized views? 3399
A.6.6 Can you insert into views that are based on joins? 3399

A.6.1. Where can I find documentation covering MySQL Views?

See Section 23.5, “Using Views”.

A.6.2. Is there a discussion forum for MySQL Views?

Yes. See http://forums.mysql.com/list.php?100

A.6.3. What happens to a view if an underlying table is dropped or renamed?

After a view has been created, it is possible to drop or alter a table or view to which the definition refers. To check a view definition for problems of this kind, use the `CHECK TABLE` statement. (See Section 13.7.2.2, “CHECK TABLE Syntax”.)

A.6.4. Does MySQL 8.0 have table snapshots?

No.

A.6.5. Does MySQL 8.0 have materialized views?

No.

A.6.6. Can you insert into views that are based on joins?

It is possible, provided that your `INSERT` statement has a column list that makes it clear there is only one table involved.

You *cannot* insert into multiple tables with a single insert on a view.

A.7 MySQL 8.0 FAQ: INFORMATION_SCHEMA

A.7.1 Where can I find documentation for the MySQL `INFORMATION_SCHEMA` database? 3399
A.7.2 Is there a discussion forum for `INFORMATION_SCHEMA`? 3399
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A.7.4 What is the difference between the Oracle Data Dictionary and MySQL `INFORMATION_SCHEMA`? 3399
A.7.5 Can I add to or otherwise modify the tables found in the `INFORMATION_SCHEMA` database? 3399

A.7.1. Where can I find documentation for the MySQL `INFORMATION_SCHEMA` database?

See Chapter 24, *INFORMATION_SCHEMA Tables*

A.7.2. Is there a discussion forum for `INFORMATION_SCHEMA`?

See http://forums.mysql.com/list.php?101.

A.7.3. Where can I find the ANSI SQL 2003 specification for `INFORMATION_SCHEMA`?

Unfortunately, the official specifications are not freely available. (ANSI makes them available for purchase.) However, there are books available, such as *SQL-99 Complete, Really* by Peter Gulutzan and Trudy Pelzer, that provide a comprehensive overview of the standard, including `INFORMATION_SCHEMA`.

A.7.4. What is the difference between the Oracle Data Dictionary and MySQL `INFORMATION_SCHEMA`?

Both Oracle and MySQL provide metadata in tables. However, Oracle and MySQL use different table names and column names. The MySQL implementation is more similar to those found in DB2 and SQL Server, which also support `INFORMATION_SCHEMA` as defined in the SQL standard.

A.7.5. Can I add to or otherwise modify the tables found in the `INFORMATION_SCHEMA` database?

No. Since applications may rely on a certain standard structure, this should not be modified. For this reason, *we cannot support bugs or other issues which result from modifying `INFORMATION_SCHEMA` tables or data*.

A.8 MySQL 8.0 FAQ: Migration

A.8.1 Where can I find information on how to migrate from MySQL 5.7 to MySQL 8.0? 3400

A.8.2 How has storage engine (table type) support changed in MySQL 8.0 from previous versions? 3400

A.8.1. Where can I find information on how to migrate from MySQL 5.7 to MySQL 8.0?

For detailed upgrade information, see Section 2.10.1, “Upgrading MySQL”. Do not skip a major version when upgrading, but rather complete the process in steps, upgrading from one major version to the next in each step. This may seem more complicated, but it will you save time and trouble. If you encounter problems during the upgrade, their origin will be easier to identify, either by you or, if you have a MySQL Enterprise subscription, by MySQL support.

A.8.2. How has storage engine (table type) support changed in MySQL 8.0 from previous versions?

Storage engine support has changed as follows:

- Support for `ISAM` tables was removed in MySQL 5.0 and you should now use the `MyISAM` storage engine in place of `ISAM`. To convert a table *`tblname`* from `ISAM` to `MyISAM`, simply issue a statement such as this one:

```
ALTER TABLE tblname ENGINE=MYISAM;
```

- Internal `RAID` for `MyISAM` tables was also removed in MySQL 5.0. This was formerly used to allow large tables in file systems that did not support file sizes greater than 2GB. All modern file systems allow for larger tables; in addition, there are now other solutions such as `MERGE` tables and views.
- The `VARCHAR` column type now retains trailing spaces in all storage engines.
- `MEMORY` tables (formerly known as `HEAP` tables) can also contain `VARCHAR` columns.

A.9 MySQL 8.0 FAQ: Security

A.9.1 Where can I find documentation that addresses security issues for MySQL? 3400
A.9.2 Does MySQL 8.0 have native support for SSL? 3401
A.9.3 Is SSL support built into MySQL binaries, or must I recompile the binary myself to enable it? 3401
A.9.4 Does MySQL 8.0 have built-in authentication against LDAP directories? 3401
A.9.5 Does MySQL 8.0 include support for Roles Based Access Control (RBAC)? 3401

A.9.1. Where can I find documentation that addresses security issues for MySQL?

The best place to start is Chapter 6, *Security*.

Other portions of the MySQL Documentation which you may find useful with regard to specific security concerns include the following:

- Section 6.1.1, “Security Guidelines”.
- Section 6.1.3, “Making MySQL Secure Against Attackers”.
- Section B.5.3.2, “How to Reset the Root Password”.
- Section 6.1.5, “How to Run MySQL as a Normal User”.
- Section 28.4.2.6, “UDF Security Precautions”.
- Section 6.1.4, “Security-Related mysqld Options and Variables”.
- Section 6.1.6, “Security Issues with LOAD DATA LOCAL”.
- Section 2.9, “Postinstallation Setup and Testing”.
- Section 6.4, “Using Secure Connections”.

A.9.2. Does MySQL 8.0 have native support for SSL?

Most 8.0 binaries have support for SSL connections between the client and server. See Section 6.4, “Using Secure Connections”.

You can also tunnel a connection using SSH, if (for example) the client application does not support SSL connections. For an example, see Section 6.4.7, “Connecting to MySQL Remotely from Windows with SSH”.

A.9.3. Is SSL support built into MySQL binaries, or must I recompile the binary myself to enable it?

Most 8.0 binaries have SSL enabled for client/server connections that are secured, authenticated, or both. See Section 6.4, “Using Secure Connections”.

A.9.4. Does MySQL 8.0 have built-in authentication against LDAP directories?

The Enterprise edition includes a PAM Authentication Plugin that supports authentication against an LDAP directory.

A.9.5. Does MySQL 8.0 include support for Roles Based Access Control (RBAC)?

Not at this time.

A.10 MySQL 8.0 FAQ: MySQL Cluster

In the following section, we answer questions that are frequently asked about MySQL Cluster and the `NDB` storage engine.

A.10.1 Which versions of the MySQL software support NDB Cluster? Do I have to compile from source? 3401
A.10.2 What do “NDB” and “NDBCLUSTER” mean? 3401
A.10.3 How many computers do I need to run an NDB Cluster, and why? 3401

A.10.1. Which versions of the MySQL software support NDB Cluster? Do I have to compile from source?

NDB Cluster is not supported in MySQL Server 8.0 releases. Instead, you should use NDB Cluster 7.5 for any new deployments; if you are using an older version of NDB Cluster, you should upgrade to this version soon as possible. For an overview of improvements made in NDB Cluster 7.5, see What is New in MySQL NDB Cluster 7.5. For an overview of improvements made in NDB Cluster 7.4, see What is New in NDB Cluster 7.4; for information about improvements made in NDB Cluster 7.3, see What is New in NDB Cluster 7.3. For information about improvements made in NDB Cluster 7.2, see What is New in MySQL NDB Cluster 7.2.

For detailed information about deploying and using NDB Cluster, see MySQL NDB Cluster 7.5 and NDB Cluster 7.6.

A.10.2. What do “NDB” and “NDBCLUSTER” mean?

“NDB” stands for “**N**etwork **D**ata**b**ase”. `NDB` and `NDBCLUSTER` are both names for the storage engine that enables clustering support with MySQL. `NDB` is preferred, but either name is correct.

A.10.3. How many computers do I need to run an NDB Cluster, and why?

A minimum of three computers is required to run a viable cluster. However, the minimum *recommended* number of computers in an NDB Cluster is four: one each to run the management and SQL nodes, and two computers to serve as data nodes. The purpose of the two data nodes is to provide redundancy; the management node must run on a separate machine to guarantee continued arbitration services in the event that one of the data nodes fails.

To provide increased throughput and high availability, you should use multiple SQL nodes (MySQL Servers connected to the cluster). It is also possible (although not strictly necessary) to run multiple management servers.

A.11 MySQL 8.0 FAQ: MySQL Chinese, Japanese, and Korean Character Sets

This set of Frequently Asked Questions derives from the experience of MySQL's Support and Development groups in handling many inquiries about CJK (Chinese-Japanese-Korean) issues.

A.11.1 What CJK character sets are available in MySQL? 3402
A.11.2 I have inserted CJK characters into my table. Why does SELECT display them as “?” characters? 3403
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A.11.4 Why do Japanese character set conversions fail? 3405
A.11.5 What should I do if I want to convert SJIS 81CA to cp932? 3406
A.11.6 How does MySQL represent the Yen (¥) sign? 3406
A.11.7 Of what issues should I be aware when working with Korean character sets in MySQL? 3406
A.11.8 Why do I get Incorrect string value error messages? 3407
A.11.9 Why does my GUI front end or browser display CJK characters incorrectly in my application using Access, PHP, or another API? 3407
A.11.10 I've upgraded to MySQL 8.0. How can I revert to behavior like that in MySQL 4.0 with regard to character sets? 3408
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A.11.15 Why are my supplementary characters rejected by MySQL? 3411
A.11.16 Should “CJK” be “CJKV”? 3411
A.11.17 Does MySQL permit CJK characters to be used in database and table names? 3412
A.11.18 Where can I find translations of the MySQL Manual into Chinese, Japanese, and Korean? 3412
A.11.19 Where can I get help with CJK and related issues in MySQL? 3412

A.11.1 What CJK character sets are available in MySQL?

The list of CJK character sets may vary depending on your MySQL version. For example, the gb18030 character set is not supported prior to MySQL 5.7.4. However, since the name of the applicable language appears in the DESCRIPTION column for every entry in the INFORMATION_SCHEMA.CHARACTER_SETS table, you can obtain a current list of all the non-Unicode CJK character sets using this query:

```
mysql> SELECT CHARACTER_SET_NAME, DESCRIPTION
       FROM INFORMATION_SCHEMA.CHARACTER_SETS
       WHERE DESCRIPTION LIKE '%Chin%'
       OR DESCRIPTION LIKE '%Japanese%'
       OR DESCRIPTION LIKE '%Korean%'
       ORDER BY CHARACTER_SET_NAME;
+--------------------+---------------------------------+
| CHARACTER_SET_NAME | DESCRIPTION                     |
+--------------------+---------------------------------+
| big5               | Big5 Traditional Chinese        |
| cp932              | SJIS for Windows Japanese       |
| eucjpms            | UJIS for Windows Japanese       |
| euckr              | EUC-KR Korean                   |
| gb18030            | China National Standard GB18030 |
| gb2312             | GB2312 Simplified Chinese       |
| gbk                | GBK Simplified Chinese          |
| sjis               | Shift-JIS Japanese              |
```

```
| ujis               | EUC-JP Japanese                  |
+--------------------+----------------------------------+
```

(For more information, see Section 24.1, “The INFORMATION_SCHEMA CHARACTER_SETS Table”.)

MySQL supports three variants of the *GB* (*Guojia Biaozhun*, or *National Standard*, or *Simplified Chinese*) character sets which are official in the People's Republic of China: `gb2312`, `gbk`, and (as of MySQL 5.7.4) `gb18030`.

Sometimes people try to insert `gbk` characters into `gb2312`, and it works most of the time because `gbk` is a superset of `gb2312`. But eventually they try to insert a rarer Chinese character and it does not work. (For an example, see Bug #16072).

Here, we try to clarify exactly what characters are legitimate in `gb2312` or `gbk`, with reference to the official documents. Please check these references before reporting `gb2312` or `gbk` bugs:

- The MySQL `gbk` character set is in reality “Microsoft code page 936”. This differs from the official `gbk` for characters `A1A4` (middle dot), `A1AA` (em dash), `A6E0-A6F5`, and `A8BB-A8C0`.
- For a listing of `gbk`/Unicode mappings, see http://www.unicode.org/Public/MAPPINGS/VENDORS/MICSFT/WINDOWS/CP936.TXT.

It is also possible to store CJK characters in Unicode character sets, although the available collations may not sort characters quite as you expect:

- The `utf8` and `ucs2` character sets support the characters from Unicode Basic Multilingual Plane (BMP). These characters have code point values between `U+0000` and `U+FFFF`.
- The `utf8mb4`, `utf16`, `utf16le`, and `utf32` character sets support BMP characters, as well as supplementary characters that lie outside the BMP. Supplementary characters have code point values between `U+10000` and `U+10FFFF`.

The collation used for a Unicode character set determines the ability to sort (that is, distinguish) characters in the set:

- Collations based on Unicode Collation Algorithm (UCA) 4.0.0 distinguish only BMP characters.
- Collations based on UCA 5.2.0 or 9.0.0 distinguish BMP and supplementary characters.
- Non-UCA collations may not distinguish all Unicode characters. For example, the `utf8mb4` default collation is `utf8mb4_general_ci`, which distinguishes only BMP characters.

Moreover, distinguishing characters is not the same as ordering them per the conventions of a given CJK language. Currently, MySQL has only one CJK-specific UCA collation, `gb18030_unicode_520_ci` (which requires use of the non-Unicode `gb18030` character set).

For information about Unicode collations and their differentiating properties, including collation properties for supplementary characters, see Section 10.1.10.1, “Unicode Character Sets”.

A.11.2. I have inserted CJK characters into my table. Why does `SELECT` display them as “?” characters?

This problem is usually due to a setting in MySQL that does not match the settings for the application program or the operating system. Here are some common steps for correcting these types of issues:

- *Be certain of what MySQL version you are using*.

  Use the statement `SELECT VERSION();` to determine this.

- *Make sure that the database is actually using the desired character set*.

People often think that the client character set is always the same as either the server character set or the character set used for display purposes. However, both of these are false assumptions. You can make sure by checking the result of `SHOW CREATE TABLE tablename` or, better yet, by using this statement:

```
SELECT character_set_name, collation_name
    FROM information_schema.columns
    WHERE table_schema = your_database_name
        AND table_name = your_table_name
        AND column_name = your_column_name;
```

- *Determine the hexadecimal value of the character or characters that are not being displayed correctly*.

  You can obtain this information for a column `column_name` in the table `table_name` using the following query:

```
SELECT HEX(column_name)
FROM table_name;
```

  `3F` is the encoding for the `?` character; this means that `?` is the character actually stored in the column. This most often happens because of a problem converting a particular character from your client character set to the target character set.

- *Make sure that a round trip is possible. When you select `literal` (or `_introducer hexadecimal-value`), do you obtain `literal` as a result*?

  For example, the Japanese Katakana character *Pe* (`ペ`) exists in all CJK character sets, and has the code point value (hexadecimal coding) `0x30da`. To test a round trip for this character, use this query:

```
SELECT 'ペ' AS `ペ`;         /* or SELECT _ucs2 0x30da; */
```

  If the result is not also ペ, the round trip failed.

  For bug reports regarding such failures, we might ask you to follow up with `SELECT HEX('ペ');`. Then we can determine whether the client encoding is correct.

- *Make sure that the problem is not with the browser or other application, rather than with MySQL*.

  Use the `mysql` client program to accomplish this task. If `mysql` displays characters correctly but your application does not, your problem is probably due to system settings.

  To determine your settings, use the `SHOW VARIABLES` statement, whose output should resemble what is shown here:

```
mysql> SHOW VARIABLES LIKE 'char%';
+--------------------------+----------------------------------------+
| Variable_name            | Value                                  |
+--------------------------+----------------------------------------+
| character_set_client     | utf8                                   |
| character_set_connection | utf8                                   |
| character_set_database   | latin1                                 |
| character_set_filesystem | binary                                 |
| character_set_results    | utf8                                   |
| character_set_server     | latin1                                 |
| character_set_system     | utf8                                   |
| character_sets_dir       | /usr/local/mysql/share/mysql/charsets/ |
+--------------------------+----------------------------------------+
```

These are typical character-set settings for an international-oriented client (notice the use of `utf8` Unicode) connected to a server in the West (`latin1` is a West Europe character set).

Although Unicode (usually the `utf8` variant on Unix, and the `ucs2` variant on Windows) is preferable to Latin, it is often not what your operating system utilities support best. Many Windows users find that a Microsoft character set, such as `cp932` for Japanese Windows, is suitable.

If you cannot control the server settings, and you have no idea what setting your underlying computer uses, try changing to a common character set for the country that you're in (`euckr` = Korea; `gb18030`, `gb2312` or `gbk` = People's Republic of China; `big5` = Taiwan; `sjis`, `ujis`, `cp932`, or `eucjpms` = Japan; `ucs2` or `utf8` = anywhere). Usually it is necessary to change only the client and connection and results settings. The `SET NAMES`. statement changes all three at once. For example:

```
SET NAMES 'big5';
```

Once the setting is correct, you can make it permanent by editing `my.cnf` or `my.ini`. For example you might add lines looking like these:

```
[mysqld]
character-set-server=big5
[client]
default-character-set=big5
```

It is also possible that there are issues with the API configuration setting being used in your application; see *Why does my GUI front end or browser not display CJK characters correctly...?* for more information.

A.11.3. What problems should I be aware of when working with the Big5 Chinese character set?

MySQL supports the Big5 character set which is common in Hong Kong and Taiwan (Republic of China). The MySQL `big5` character set is in reality Microsoft code page 950, which is very similar to the original `big5` character set.

A feature request for adding `HKSCS` extensions has been filed. People who need this extension may find the suggested patch for Bug #13577 to be of interest.

A.11.4. Why do Japanese character set conversions fail?

MySQL supports the `sjis`, `ujis`, `cp932`, and `eucjpms` character sets, as well as Unicode. A common need is to convert between character sets. For example, there might be a Unix server (typically with `sjis` or `ujis`) and a Windows client (typically with `cp932`).

In the following conversion table, the `ucs2` column represents the source, and the `sjis`, `cp932`, `ujis`, and `eucjpms` columns represent the destinations; that is, the last 4 columns provide the hexadecimal result when we use `CONVERT(ucs2)` or we assign a `ucs2` column containing the value to an `sjis`, `cp932`, `ujis`, or `eucjpms` column.

Character Name	ucs2	sjis	cp932	ujis	eucjpms
BROKEN BAR	00A6	3F	3F	8FA2C3	3F
FULLWIDTH BROKEN BAR	FFE4	3F	FA55	3F	8FA2
YEN SIGN	00A5	3F	3F	20	3F
FULLWIDTH YEN SIGN	FFE5	818F	818F	A1EF	3F
TILDE	007E	7E	7E	7E	7E
OVERLINE	203E	3F	3F	20	3F

Character Name	ucs2	sjis	cp932	ujis	eucjpms
HORIZONTAL BAR	2015	815C	815C	A1BD	A1BD
EM DASH	2014	3F	3F	3F	3F
REVERSE SOLIDUS	005C	815F	5C	5C	5C
FULLWIDTH ""	FF3C	3F	815F	3F	A1C0
WAVE DASH	301C	8160	3F	A1C1	3F
FULLWIDTH TILDE	FF5E	3F	8160	3F	A1C1
DOUBLE VERTICAL LINE	2016	8161	3F	A1C2	3F
PARALLEL TO	2225	3F	8161	3F	A1C2
MINUS SIGN	2212	817C	3F	A1DD	3F
FULLWIDTH HYPHEN-MINUS	FF0D	3F	817C	3F	A1DD
CENT SIGN	00A2	8191	3F	A1F1	3F
FULLWIDTH CENT SIGN	FFE0	3F	8191	3F	A1F1
POUND SIGN	00A3	8192	3F	A1F2	3F
FULLWIDTH POUND SIGN	FFE1	3F	8192	3F	A1F2
NOT SIGN	00AC	81CA	3F	A2CC	3F
FULLWIDTH NOT SIGN	FFE2	3F	81CA	3F	A2CC

Now consider the following portion of the table.

	ucs2	sjis	cp932
NOT SIGN	00AC	81CA	3F
FULLWIDTH NOT SIGN	FFE2	3F	81CA

This means that MySQL converts the `NOT SIGN` (Unicode `U+00AC`) to `sjis` code point `0x81CA` and to `cp932` code point `3F`. (`3F` is the question mark (“?”. This is what is always used when the conversion cannot be performed.)

A.11.5 What should I do if I want to convert SJIS `81CA` to `cp932`?

Our answer is: “?”. There are disadvantages to this, and many people would prefer a “loose” conversion, so that `81CA (NOT SIGN)` in `sjis` becomes `81CA (FULLWIDTH NOT SIGN)` in `cp932`.

A.11.6 How does MySQL represent the Yen (`¥`) sign?

A problem arises because some versions of Japanese character sets (both `sjis` and `euc`) treat `5C` as a *reverse solidus* (`\`, also known as a backslash), whereas others treat it as a yen sign (`¥`).

MySQL follows only one version of the JIS (Japanese Industrial Standards) standard description. In MySQL, *`5C` is always the reverse solidus (`\`)*.

A.11.7 Of what issues should I be aware when working with Korean character sets in MySQL?

In theory, while there have been several versions of the `euckr` (*Extended Unix Code Korea*) character set, only one problem has been noted. We use the “ASCII” variant of EUC-KR, in which the code point `0x5c` is REVERSE SOLIDUS, that is `\`, instead of the “KS-Roman” variant of EUC-KR, in which the code point `0x5c` is `WON SIGN` (`₩`). This means that you cannot convert Unicode `U+20A9` to `euckr`:

```
mysql> SELECT
```

```
          CONVERT('₩' USING euckr) AS euckr,
          HEX(CONVERT('₩' USING euckr)) AS hexeuckr;
+-------+----------+
| euckr | hexeuckr |
+-------+----------+
| ?     | 3F       |
+-------+----------+
```

A.11.8. Why do I get `Incorrect string value` error messages?

To see the problem, create a table with one Unicode (`ucs2`) column and one Chinese (`gb2312`) column.

```
mysql> CREATE TABLE ch
       (ucs2 CHAR(3) CHARACTER SET ucs2,
       gb2312 CHAR(3) CHARACTER SET gb2312);
```

In nonstrict SQL mode, try to place the rare character 浻 in both columns.

```
mysql> SET sql_mode = '';
mysql> INSERT INTO ch VALUES ('A浻B','A浻B');
Query OK, 1 row affected, 1 warning (0.00 sec)
```

The `INSERT` produces a warning. Use the following statement to see what it is:

```
mysql> SHOW WARNINGS\G
*************************** 1. row ***************************
  Level: Warning
   Code: 1366
Message: Incorrect string value: '\xE6\xB1\x8CB' for column 'gb2312' at row 1
```

So it is a warning about the `gb2312` column only.

```
mysql> SELECT ucs2,HEX(ucs2),gb2312,HEX(gb2312) FROM ch;
+-------+--------------+--------+-------------+
| ucs2  | HEX(ucs2)    | gb2312 | HEX(gb2312) |
+-------+--------------+--------+-------------+
| A浻B | 00416C4C0042 | A?B    | 413F42      |
+-------+--------------+--------+-------------+
```

Several things need explanation here:

1. The 浻 character is not in the `gb2312` character set, as described earlier.
2. If you are using an old version of MySQL, you may see a different message.
3. A warning occurs rather than an error because MySQL is not set to use strict SQL mode. In nonstrict mode, MySQL tries to do what it can, to get the best fit, rather than give up. With strict SQL mode, the `Incorrect string value` message occurs as an error rather than a warning, and the `INSERT` fails.

A.11.9. Why does my GUI front end or browser display CJK characters incorrectly in my application using Access, PHP, or another API?

Obtain a direct connection to the server using the `mysql` client, and try the same query there. If `mysql` responds correctly, the trouble may be that your application interface requires initialization. Use `mysql` to tell you what character set or sets it uses with the statement `SHOW VARIABLES LIKE 'char%';`. If you are using Access, you are most likely connecting with Connector/ODBC. In this case, you should check Configuring Connector/ODBC. If, for example, you use `big5`, you would enter `SET NAMES 'big5'`. (In this case, no `;` character is required.) If you are using ASP, you might need to add `SET NAMES` in the code. Here is an example that has worked in the past:

```
<%
Session.CodePage=0
Dim strConnection
Dim Conn
strConnection="driver={MySQL ODBC 3.51 Driver};server=server;uid=username;" \
               & "pwd=password;database=database;stmt=SET NAMES 'big5';"
Set Conn = Server.CreateObject("ADODB.Connection")
Conn.Open strConnection
%>
```

In much the same way, if you are using any character set other than `latin1` with Connector/Net, you must specify the character set in the connection string. See Connecting to MySQL Using Connector/Net, for more information.

If you are using PHP, try this:

```
<?php
  $link = new mysqli($host, $usr, $pwd, $db);

  if( mysqli_connect_errno() )
  {
    printf("Connect failed: %s\n", mysqli_connect_error());
    exit();
  }

  $link->query("SET NAMES 'utf8'");
?>
```

In this case, we used `SET NAMES` to change `character_set_client`, `character_set_connection`, and `character_set_results`.

Another issue often encountered in PHP applications has to do with assumptions made by the browser. Sometimes adding or changing a `<meta>` tag suffices to correct the problem: for example, to insure that the user agent interprets page content as `UTF-8`, include `<meta http-equiv="Content-Type" content="text/html; charset=utf-8">` in the `<head>` section of the HTML page.

If you are using Connector/J, see Using Character Sets and Unicode.

A.11.10 ve upgraded to MySQL 8.0. How can I revert to behavior like that in MySQL 4.0 with regard to character sets?

In MySQL Version 4.0, there was a single “global” character set for both server and client, and the decision as to which character to use was made by the server administrator. This changed starting with MySQL Version 4.1. What happens now is a “handshake”, as described in Section 10.1.4, “Connection Character Sets and Collations”:

> When a client connects, it sends to the server the name of the character set that it wants to use. The server uses the name to set the `character_set_client`, `character_set_results`, and `character_set_connection` system variables. In effect, the server performs a `SET NAMES` operation using the character set name.

The effect of this is that you cannot control the client character set by starting `mysqld` with `--character-set-server=utf8`. However, some Asian customers prefer the MySQL 4.0 behavior. To make it possible to retain this behavior, we added a `mysqld` switch, `--character-set-client-handshake`, which can be turned off with `--skip-character-set-client-handshake`. If you start `mysqld` with `--skip-character-set-client-handshake`, then, when a client connects, it sends to the server the name of the character set that it wants to use. However, *the server ignores this request from the client*.

By way of example, suppose that your favorite server character set is `latin1` (unlikely in a CJK area, but this is the default value). Suppose further that the client uses `utf8` because this is what the client's operating system supports. Now, start the server with `latin1` as its default character set:

```
mysqld --character-set-server=latin1
```

And then start the client with the default character set `utf8`:

```
mysql --default-character-set=utf8
```

The resulting settings can be seen by viewing the output of `SHOW VARIABLES`:

```
mysql> SHOW VARIABLES LIKE 'char%';
+--------------------------+----------------------------------------+
| Variable_name            | Value                                  |
+--------------------------+----------------------------------------+
| character_set_client     | utf8                                   |
| character_set_connection | utf8                                   |
| character_set_database   | latin1                                 |
| character_set_filesystem | binary                                 |
| character_set_results    | utf8                                   |
| character_set_server     | latin1                                 |
| character_set_system     | utf8                                   |
| character_sets_dir       | /usr/local/mysql/share/mysql/charsets/ |
+--------------------------+----------------------------------------+
```

Now stop the client, and stop the server using `mysqladmin`. Then start the server again, but this time tell it to skip the handshake like so:

```
mysqld --character-set-server=utf8 --skip-character-set-client-handshake
```

Start the client with `utf8` once again as the default character set, then display the resulting settings:

```
mysql> SHOW VARIABLES LIKE 'char%';
+--------------------------+----------------------------------------+
| Variable_name            | Value                                  |
+--------------------------+----------------------------------------+
| character_set_client     | latin1                                 |
| character_set_connection | latin1                                 |
| character_set_database   | latin1                                 |
| character_set_filesystem | binary                                 |
| character_set_results    | latin1                                 |
| character_set_server     | latin1                                 |
| character_set_system     | utf8                                   |
| character_sets_dir       | /usr/local/mysql/share/mysql/charsets/ |
+--------------------------+----------------------------------------+
```

As you can see by comparing the differing results from `SHOW VARIABLES`, the server ignores the client's initial settings if the `--skip-character-set-client-handshake` option is used.

A.11.11 Why do some `LIKE` and `FULLTEXT` searches with CJK characters fail?

For `LIKE` searches, there is a very simple problem with binary string column types such as `BINARY` and `BLOB`: we must know where characters end. With multibyte character sets, different characters might have different octet lengths. For example, in `utf8`, `A` requires one byte but ペ requires three bytes, as shown here:

```
+-------------------------+--------------------------+
| OCTET_LENGTH(_utf8 'A') | OCTET_LENGTH(_utf8 'ペ') |
```

```
+------------------------+--------------------------+
|                      1 |                        3 |
+------------------------+--------------------------+
```

If we do not know where the first character in a string ends, we do not know where the second character begins, in which case even very simple searches such as `LIKE '_A%'` fail. The solution is to use a nonbinary string column type defined to have the proper CJK character set. For example: `mycol TEXT CHARACTER SET sjis`. Alternatively, convert to a CJK character set before comparing.

This is one reason why MySQL cannot permit encodings of nonexistent characters. If it is not strict about rejecting bad input, it has no way of knowing where characters end.

For `FULLTEXT` searches, we must know where words begin and end. With Western languages, this is rarely a problem because most (if not all) of these use an easy-to-identify word boundary: the space character. However, this is not usually the case with Asian writing. We could use arbitrary halfway measures, like assuming that all Han characters represent words, or (for Japanese) depending on changes from Katakana to Hiragana due to grammatical endings. However, the only sure solution requires a comprehensive word list, which means that we would have to include a dictionary in the server for each Asian language supported. This is simply not feasible.

A.11.12. How do I know whether character *X* is available in all character sets?

The majority of simplified Chinese and basic nonhalfwidth Japanese Kana characters appear in all CJK character sets. The following stored procedure accepts a `UCS-2` Unicode character, converts it to other character sets, and displays the results in hexadecimal.

```
DELIMITER //

CREATE PROCEDURE p_convert(ucs2_char CHAR(1) CHARACTER SET ucs2)
BEGIN

CREATE TABLE tj
             (ucs2 CHAR(1) character set ucs2,
              utf8 CHAR(1) character set utf8,
              big5 CHAR(1) character set big5,
              cp932 CHAR(1) character set cp932,
              eucjpms CHAR(1) character set eucjpms,
              euckr CHAR(1) character set euckr,
              gb2312 CHAR(1) character set gb2312,
              gbk CHAR(1) character set gbk,
              sjis CHAR(1) character set sjis,
              ujis CHAR(1) character set ujis);

INSERT INTO tj (ucs2) VALUES (ucs2_char);

UPDATE tj SET utf8=ucs2,
              big5=ucs2,
              cp932=ucs2,
              eucjpms=ucs2,
              euckr=ucs2,
              gb2312=ucs2,
              gbk=ucs2,
              sjis=ucs2,
              ujis=ucs2;

/* If there are conversion problems, UPDATE produces warnings. */

SELECT hex(ucs2) AS ucs2,
       hex(utf8) AS utf8,
       hex(big5) AS big5,
       hex(cp932) AS cp932,
       hex(eucjpms) AS eucjpms,
       hex(euckr) AS euckr,
       hex(gb2312) AS gb2312,
       hex(gbk) AS gbk,
```

```
      hex(sjis) AS sjis,
      hex(ujis) AS ujis
FROM tj;

DROP TABLE tj;

END//

DELIMITER ;
```

The input can be any single `ucs2` character, or it can be the code value (hexadecimal representation) of that character. For example, from Unicode's list of `ucs2` encodings and names (http://www.unicode.org/Public/UNIDATA/UnicodeData.txt), we know that the Katakana character *Pe* appears in all CJK character sets, and that its code value is `X'30DA'`. If we use this value as the argument to `p_convert()`, the result is as shown here:

```
mysql> CALL p_convert(X'30DA');
+------+--------+------+-------+---------+-------+--------+------+------+------+
| ucs2 | utf8   | big5 | cp932 | eucjpms | euckr | gb2312 | gbk  | sjis | ujis |
+------+--------+------+-------+---------+-------+--------+------+------+------+
| 30DA | E3839A | C772 | 8379  | A5DA    | ABDA  | A5DA   | A5DA | 8379 | A5DA |
+------+--------+------+-------+---------+-------+--------+------+------+------+
```

Since none of the column values is `3F` (that is, the question mark character, `?`), we know that every conversion worked.

A.11.13. Why do CJK strings sort incorrectly in Unicode? (I)

CJK sorting problems that occurred in older MySQL versions can be solved as of MySQL 8.0 by using the `utf8mb4` character set and the `utf8mb4_ja_0900_as_cs` collation.

A.11.14. Why do CJK strings sort incorrectly in Unicode? (II)

CJK sorting problems that occurred in older MySQL versions can be solved as of MySQL 8.0 by using the `utf8mb4` character set and the `utf8mb4_ja_0900_as_cs` collation.

A.11.15. Why are my supplementary characters rejected by MySQL?

Supplementary characters lie outside the Unicode *Basic Multilingual Plane / Plane 0*. BMP characters have code point values between `U+0000` and `U+FFFF`. Supplementary characters have code point values between `U+10000` and `U+10FFFF`.

To store supplementary characters, you must use a character set that permits them:

- The `utf8` and `ucs2` character sets support BMP characters only.

  The `utf8` character set permits only `UTF-8` characters that take up to three bytes. This has led to reports such as that found in Bug #12600, which we rejected as “not a bug”. With `utf8`, MySQL must truncate an input string when it encounters bytes that it does no understand. Otherwise, it is unknown how long the bad multibyte character is.

  One possible workaround is to use `ucs2` instead of `utf8`, in which case the “bad” characters are changed to question marks. However, no truncation takes place. You can also change the data type to `BLOB` or `BINARY`, which perform no validity checking.

- The `utf8mb4`, `utf16`, `utf16le`, and `utf32` character sets support BMP characters, as well as supplementary characters outside the BMP.

A.11.16. Should “CJK” be “CJKV”?

No. The term “CJKV” (*Chinese Japanese Korean Vietnamese*) refers to Vietnamese character sets which contain Han (originally Chinese) characters. MySQL supports the modern Vietnamese script with Western characters, but does not support the old Vietnamese script using Han characters.

As of MySQL 5.6, there are Vietnamese collations for Unicode character sets, as described in Section 10.1.10.1, “Unicode Character Sets”.

A.11.17 Does MySQL permit CJK characters to be used in database and table names?

Yes.

A.11.18 Where can I find translations of the MySQL Manual into Chinese, Japanese, and Korean?

The Japanese translation of the MySQL 5.6 manual can be downloaded from http://dev.mysql.com/doc/.

A.11.19 Where can I get help with CJK and related issues in MySQL?

The following resources are available:

- A listing of MySQL user groups can be found at https://wikis.oracle.com/display/mysql/List+of+MySQL+User+Groups.
- View feature requests relating to character set issues at http://tinyurl.com/y6xcuf.
- Visit the MySQL Character Sets, Collation, Unicode Forum. http://forums.mysql.com/ also provides foreign-language forums.

A.12 MySQL 8.0 FAQ: Connectors & APIs

For common questions, issues, and answers relating to the MySQL Connectors and other APIs, see the following areas of the Manual:

- Section 27.7.18, “Common Questions and Problems When Using the C API”
- Common Problems with MySQL and PHP
- Connector/ODBC Notes and Tips
- Connector/Net Programming
- MySQL Connector/J 5.1 Developer Guide

A.13 MySQL 8.0 FAQ: Replication

In the following section, we provide answers to questions that are most frequently asked about MySQL Replication.

A.13.1 Must the slave be connected to the master all the time? 3413
A.13.2 Must I enable networking on my master and slave to enable replication? 3413
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A.13.14 Does replication work on mixed hardware architectures (for example, the master runs on a 64-bit machine while slaves run on 32-bit machines)? 3416

A.13.1 Must the slave be connected to the master all the time?

No, it does not. The slave can go down or stay disconnected for hours or even days, and then reconnect and catch up on updates. For example, you can set up a master/slave relationship over a dial-up link where the link is up only sporadically and for short periods of time. The implication of this is that, at any given time, the slave is not guaranteed to be in synchrony with the master unless you take some special measures.

To ensure that catchup can occur for a slave that has been disconnected, you must not remove binary log files from the master that contain information that has not yet been replicated to the slaves. Asynchronous replication can work only if the slave is able to continue reading the binary log from the point where it last read events.

A.13.2 Must I enable networking on my master and slave to enable replication?

Yes, networking must be enabled on the master and slave. If networking is not enabled, the slave cannot connect to the master and transfer the binary log. Check that the `skip-networking` option has not been enabled in the configuration file for either server.

A.13.3 How do I know how late a slave is compared to the master? In other words, how do I know the date of the last statement replicated by the slave?

Check the `Seconds_Behind_Master` column in the output from `SHOW SLAVE STATUS`. See Section 18.1.7.1, “Checking Replication Status”.

When the slave SQL thread executes an event read from the master, it modifies its own time to the event timestamp. (This is why `TIMESTAMP` is well replicated.) In the `Time` column in the output of `SHOW PROCESSLIST`, the number of seconds displayed for the slave SQL thread is the number of seconds between the timestamp of the last replicated event and the real time of the slave machine. You can use this to determine the date of the last replicated event. Note that if your slave has been disconnected from the master for one hour, and then reconnects, you may immediately see large `Time` values such as 3600 for the slave SQL thread in `SHOW PROCESSLIST`. This is because the slave is executing statements that are one hour old. See Section 18.2.2, “Replication Implementation Details”.

A.13.4 How do I force the master to block updates until the slave catches up?

Use the following procedure:

1. On the master, execute these statements:

```
mysql> FLUSH TABLES WITH READ LOCK;
mysql> SHOW MASTER STATUS;
```

   Record the replication coordinates (the current binary log file name and position) from the output of the `SHOW` statement.

2. On the slave, issue the following statement, where the arguments to the `MASTER_POS_WAIT()` function are the replication coordinate values obtained in the previous step:

```
mysql> SELECT MASTER_POS_WAIT('log_name', log_pos);
```

   The `SELECT` statement blocks until the slave reaches the specified log file and position. At that point, the slave is in synchrony with the master and the statement returns.

3. On the master, issue the following statement to enable the master to begin processing updates again:

```
mysql> UNLOCK TABLES;
```

A.13.5 What issues should I be aware of when setting up two-way replication?

MySQL replication currently does not support any locking protocol between master and slave to guarantee the atomicity of a distributed (cross-server) update. In other words, it is possible for client A to make an update to co-master 1, and in the meantime, before it propagates to co-master 2, client B could make an update to co-master 2 that makes the update of client A work differently than it did on co-master 1. Thus, when the update of client A makes it to co-master 2, it produces tables that are different from what you have on co-master 1, even after all the updates from co-master 2 have also propagated. This means that you should not chain two servers together in a two-way replication relationship unless you are sure that your updates can safely happen in any order, or unless you take care of mis-ordered updates somehow in the client code.

You should also realize that two-way replication actually does not improve performance very much (if at all) as far as updates are concerned. Each server must do the same number of updates, just as you would have a single server do. The only difference is that there is a little less lock contention because the updates originating on another server are serialized in one slave thread. Even this benefit might be offset by network delays.

A.13.6 How can I use replication to improve performance of my system?

Set up one server as the master and direct all writes to it. Then configure as many slaves as you have the budget and rackspace for, and distribute the reads among the master and the slaves. You can also start the slaves with the `--skip-innodb`, `--low-priority-updates`, and `--delay-key-write=ALL` options to get speed improvements on the slave end. In this case, the slave uses nontransactional `MyISAM` tables instead of `InnoDB` tables to get more speed by eliminating transactional overhead.

A.13.7 What should I do to prepare client code in my own applications to use performance-enhancing replication?

See the guide to using replication as a scale-out solution, Section 18.3.5, “Using Replication for Scale-Out”.

A.13.8 When and how much can MySQL replication improve the performance of my system?

MySQL replication is most beneficial for a system that processes frequent reads and infrequent writes. In theory, by using a single-master/multiple-slave setup, you can scale the system by adding more slaves until you either run out of network bandwidth, or your update load grows to the point that the master cannot handle it.

To determine how many slaves you can use before the added benefits begin to level out, and how much you can improve performance of your site, you must know your query patterns, and determine empirically by benchmarking the relationship between the throughput for reads and writes on a typical master and a typical slave. The example here shows a rather simplified calculation of what you can get with replication for a hypothetical system. Let `reads` and `writes` denote the number of reads and writes per second, respectively.

Let's say that system load consists of 10% writes and 90% reads, and we have determined by benchmarking that `reads` is 1200 - 2 * `writes`. In other words, the system can do 1,200 reads per second with no writes, the average write is twice as slow as the average read, and the relationship is linear. Suppose that the master and each slave have the same capacity, and that we have one master and *`N`* slaves. Then we have for each server (master or slave):

`reads` = 1200 - 2 * `writes`

`reads` = 9 * `writes` / (*`N`* + 1) (reads are split, but writes replicated to all slaves)

9 * `writes` / (*N* + 1) + 2 * `writes` = 1200

`writes` = 1200 / (2 + 9/(*N* + 1))

The last equation indicates the maximum number of writes for *N* slaves, given a maximum possible read rate of 1,200 per second and a ratio of nine reads per write.

This analysis yields the following conclusions:

- If *N* = 0 (which means we have no replication), our system can handle about 1200/11 = 109 writes per second.
- If *N* = 1, we get up to 184 writes per second.
- If *N* = 8, we get up to 400 writes per second.
- If *N* = 17, we get up to 480 writes per second.
- Eventually, as *N* approaches infinity (and our budget negative infinity), we can get very close to 600 writes per second, increasing system throughput about 5.5 times. However, with only eight servers, we increase it nearly four times.

These computations assume infinite network bandwidth and neglect several other factors that could be significant on your system. In many cases, you may not be able to perform a computation similar to the one just shown that accurately predicts what will happen on your system if you add *N* replication slaves. However, answering the following questions should help you decide whether and by how much replication will improve the performance of your system:

- What is the read/write ratio on your system?
- How much more write load can one server handle if you reduce the reads?
- For how many slaves do you have bandwidth available on your network?

A.13.9. How can I use replication to provide redundancy or high availability?

How you implement redundancy is entirely dependent on your application and circumstances. High-availability solutions (with automatic failover) require active monitoring and either custom scripts or third party tools to provide the failover support from the original MySQL server to the slave.

To handle the process manually, you should be able to switch from a failed master to a pre-configured slave by altering your application to talk to the new server or by adjusting the DNS for the MySQL server from the failed server to the new server.

For more information and some example solutions, see Section 18.3.8, “Switching Masters During Failover”.

A.13.10. How do I tell whether a master server is using statement-based or row-based binary logging format?

Check the value of the `binlog_format` system variable:

```
mysql> SHOW VARIABLES LIKE 'binlog_format';
```

The value shown will be one of `STATEMENT`, `ROW`, or `MIXED`. For `MIXED` mode, statement-based logging is used by default but replication switches automatically to row-based logging under certain conditions, such as unsafe statements. For information about when this may occur, see Section 5.4.4.3, “Mixed Binary Logging Format”.

A.13.11. How do I tell a slave to use row-based replication?

Slaves automatically know which format to use.

A.13.12 How do I prevent `GRANT` and `REVOKE` statements from replicating to slave machines?

Start the server with the `--replicate-wild-ignore-table=mysql.%` option to ignore replication for tables in the `mysql` database.

A.13.13 Does replication work on mixed operating systems (for example, the master runs on Linux while slaves run on OS X and Windows)?

Yes.

A.13.14 Does replication work on mixed hardware architectures (for example, the master runs on a 64-bit machine while slaves run on 32-bit machines)?

Yes.

A.14 MySQL 8.0 FAQ: InnoDB Change Buffer

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A.14.1 What types of operations modify secondary indexes and result in change buffering?

`INSERT`, `UPDATE`, and `DELETE` operations can modify secondary indexes. If an affected index page is not in the buffer pool, the changes can be buffered in the change buffer.

A.14.2 What is the benefit of the `InnoDB` change buffer?

Buffering secondary index changes when secondary index pages are not in the buffer pool avoids expensive random access I/O operations that would be required to immediately read in affected index pages from disk. Buffered changes can be applied later, in batches, as pages are read into the buffer pool by other read operations.

A.14.3 Does the change buffer support other types of indexes?

No. The change buffer only supports secondary indexes. Clustered indexes, full-text indexes, and spatial indexes are not supported. Full-text indexes have their own caching mechanism.

A.14.4 How much space does `InnoDB` use for the change buffer?

Prior to the introduction of the `innodb_change_buffer_max_size` configuration option in MySQL 5.6, the maximum size of the on-disk change buffer in the system tablespace was 1/3 of the `InnoDB` buffer pool size.

In MySQL 5.6 and later, the `innodb_change_buffer_max_size` configuration option defines the maximum size of the change buffer as a percentage of the total buffer pool size. By default, `innodb_change_buffer_max_size` is set to 25. The maximum setting is 50.

`InnoDB` does not buffer an operation if it would cause the on-disk change buffer to exceed the defined limit.

Change buffer pages are not required to persist in the buffer pool and may be evicted by LRU operations.

A.14.5 How do I determine the current size of the change buffer?

The current size of the change buffer is reported by `SHOW ENGINE INNODB STATUS \G`, under the `INSERT BUFFER AND ADAPTIVE HASH INDEX` heading. For example:

```
-------------------------------------
INSERT BUFFER AND ADAPTIVE HASH INDEX
-------------------------------------
Ibuf: size 1, free list len 0, seg size 2, 0 merges
```

Relevant data points include:

- `size`: The number of pages used within the change buffer. Change buffer size is equal to `seg size - (1 + free list len)`. The `1 +` value represents the change buffer header page.
- `seg size`: The size of the change buffer, in pages.

For information about monitoring change buffer status, see Section 15.4.2, “Change Buffer”.

A.14.6 When does change buffer merging occur?

- When a page is read into the buffer pool, buffered changes are merged upon completion of the read, before the page is made available.
- Change buffer merging is performed as a background task. The `innodb_io_capacity` parameter sets an upper limit on the I/O activity performed by `InnoDB` background tasks such as merging data from the change buffer.
- A change buffer merge is performed during crash recovery. Changes are applied from the change buffer (in the system tablespace) to leaf pages of secondary indexes as index pages are read into the buffer pool.
- The change buffer is fully durable and will survive a system crash. Upon restart, change buffer merge operations resume as part of normal operations.
- A full merge of the change buffer can be forced as part of a slow server shutdown using `--innodb-fast-shutdown=0`.

A.14.7 When is the change buffer flushed?

Updated pages are flushed by the same flushing mechanism that flushes the other pages that occupy the buffer pool.

A.14.8 When should the change buffer be used?

The change buffer is a feature designed to reduce random I/O to secondary indexes as indexes grow larger and no longer fit in the `InnoDB` buffer pool. Generally, the change buffer should be used when the entire data set does not fit into the buffer pool, when there is substantial DML activity that modifies secondary index pages, or when there are lots of secondary indexes that are regularly changed by DML activity.

A.14.9 When should the change buffer not be used?

You might consider disabling the change buffer if the entire data set fits within the `InnoDB` buffer pool, if you have relatively few secondary indexes, or if you are using solid-state storage, where random reads are about as fast as sequential reads. Before making configuration changes, it is recommended that you run tests using a representative workload to determine if disabling the change buffer provides any benefit.

A.14.10 Where can I find additional information about the change buffer?

The following sections and blog post provide additional information about the `InnoDB` change buffer:

- Section 15.4.2, “Change Buffer”
- Section 15.6.4, “Configuring InnoDB Change Buffering”
- MySQL 5.5: InnoDB Change Buffering

A.15 MySQL 8.0 FAQ: InnoDB Tablespace Encryption

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A.15.1 Is data decrypted for users who are authorized to see it?

Yes. InnoDB tablespace encryption is designed to provide customers with the ability to transparently apply encryption within the database without impacting existing applications. Returning data in encrypted format would break most existing applications. InnoDB tablespace encryption provides the benefit of encryption without the overhead associated with traditional database encryption solutions, which would typically require expensive and substantial changes to applications, database triggers, and views.

A.15.2 What is the overhead associated with InnoDB tablespace encryption?

There is no additional storage overhead. According to internal benchmarks, performance overhead amounts to a single digit percentage difference.

A.15.3 What are the encryption algorithms used with InnoDB tablespace encryption?

InnoDB tablespace encryption supports the Advanced Encryption Standard (AES256) block-based encryption algorithm. It uses Electronic Codebook (ECB) block encryption mode for tablespace key encryption and Cipher Block Chaining (CBC) block encryption mode for data encryption.

A.15.4 Is it possible to use 3rd party encryption algorithms in place of the one provided by the InnoDB tablespace encryption feature?

No, it is not possible to use other encryption algorithms. The provided encryption algorithm is broadly accepted.

A.15.5 Can indexed columns be encrypted?

InnoDB tablespace encryption supports all indexes transparently.

A.15.6 What data types and data lengths does InnoDB tablespace encryption support?

`InnoDB` tablespace encryption supports all supported data types. There is no data length limitation.

A.15.7 Does data remain encrypted on the network?

Data encrypted by the `InnoDB` tablespace encryption feature is decrypted when it is read from the tablespace file. Thus, if the data is on the network, it is in clear-text form. However, data on the network can be encrypted using MySQL network encryption, which encrypts data traveling to and from a database using SSL/TLS.

A.15.8 Does database memory contain clear-text or encrypted data?

With `InnoDB` tablespace encryption, in-memory data is decrypted, which provides 100% transparency.

A.15.9 How do I know which data to encrypt?

Compliance with the PCI-DSS standard requires that credit card numbers (Primary Account Number, or 'PAN') be stored in encrypted form. Breach Notification Laws (for example, CA SB 1386, CA AB 1950, and similar laws in 43+ more US states) require encryption of first name, last name, driver license number, and other PII data. In early 2008, CA AB 1298 added medical and health insurance information to PII data. Additionally, industry specific privacy and security standards may require encryption of certain assets. For example, assets such as pharmaceutical research results, oil field exploration results, financial contracts, or personal data of law enforcement informants may require encryption. In the health care industry, the privacy of patient data, health records and X-ray images is of the highest importance.

A.15.10 How is `InnoDB` tablespace encryption different from encryption functions MySQL already provides?

There are symmetric and asymmetric encryption APIs in MySQL that can be used to manually encrypt data within the database. However, the application must manage encryption keys and perform required encryption and decryption operations by calling API functions. `InnoDB` tablespace encryption requires no application changes, is transparent to end users, and provides automated, built-in key management.

A.15.11 Does the transportable tablespaces feature work with `InnoDB` tablespace encryption?

Yes. For more information, see Exporting Encrypted Tables.

A.15.12 Does compression work with `InnoDB` tablespace encryption?

Customers using `InnoDB` tablespace encryption receive the full benefit of compression because compression is applied before data blocks are encrypted.

A.15.13 Can I use `mysqlpump` or `mysqldump` with encrypted tables?

Yes. Because these utilities create logical backups, the data dumped from encrypted tables is not encrypted.

A.15.14 How do I change (rotate, re-key) the master encryption key?

`InnoDB` tablespace encryption uses a two tier key mechanism. When tablespace encryption is used, individual tablespace keys are stored in the header of the underlying tablespace data file. Tablespace keys are encrypted using the master encryption key. The master encryption key is generated when tablespace encryption is enabled, and is stored outside the database. The master encryption key is rotated using the `ALTER INSTANCE ROTATE INNODB MASTER KEY` statement, which generates a new master encryption key, stores the key, and rotates the key into use.

A.15.15 How do I migrate data from a clear-text `InnoDB` tablespace to an encrypted `InnoDB` tablespace?

Transferring data from one tablespace to another is not required. To encrypt data in an `InnoDB` tablespace, run `ALTER TABLE tbl_name ENCRYPTION='Y'`. `InnoDB` tablespace encryption is only supported with file-per-table tablespaces.

A.16 MySQL 8.0 FAQ: Virtualization Support

A.16.1 Is MySQL supported on virtualized environments such as Oracle VM, VMWare, Docker, Microsoft Hyper-V, or others? 3420

A.16.1 Is MySQL supported on virtualized environments such as Oracle VM, VMWare, Docker, Microsoft Hyper-V, or others?

MySQL is supported on virtualized environments, but is certified only for Oracle VM. Contact Oracle Support for more information.

Be aware of potential problems when using virtualization software. The usual ones are related to performance, performance degradations, slowness, or unpredictability of disk, I/O, network, and memory.

Appendix B Errors, Error Codes, and Common Problems

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This appendix lists common problems and errors that may occur and potential resolutions, in addition to listing the errors that may appear when you call MySQL from any host language. The first section covers problems and resolutions. Detailed information on errors is provided: One list displays server error messages. Another list displays client program messages.

B.1 Sources of Error Information

There are several sources of error information in MySQL:

- Each SQL statement executed results in an error code, an SQLSTATE value, and an error message, as described in Section B.2, “Types of Error Values”. These errors are returned from the server side; see Section B.3, “Server Error Codes and Messages”.
- Errors can occur on the client side, usually involving problems communicating with the server; see Section B.4, “Client Error Codes and Messages”.
- SQL statement warning and error information is available through the `SHOW WARNINGS` and `SHOW ERRORS` statements. The `warning_count` system variable indicates the number of errors, warnings, and notes. The `error_count` system variable indicates the number of errors. Its value excludes warnings and notes.
- The `GET DIAGNOSTICS` statement may be used to inspect the diagnostic information in the diagnostics area. See Section 13.6.7.3, “GET DIAGNOSTICS Syntax”.
- `SHOW SLAVE STATUS` statement output includes information about replication errors occurring on the slave side.
- `SHOW ENGINE INNODB STATUS` statement output includes information about the most recent foreign key error if a `CREATE TABLE` statement for an `InnoDB` table fails.
- The `perror` program provides information from the command line about error numbers. See Section 4.8.2, “`perror` — Explain Error Codes”.

Descriptions of server and client errors are provided later in this Appendix. For information about errors related to `InnoDB`, see Section 15.20.4, “InnoDB Error Handling”.

B.2 Types of Error Values

When an error occurs in MySQL, the server returns two types of error values:

- A MySQL-specific error code. This value is numeric. It is not portable to other database systems.

- An SQLSTATE value. The value is a five-character string (for example, `'42S02'`). The values are taken from ANSI SQL and ODBC and are more standardized.

A message string that provides a textual description of the error is also available.

When an error occurs, the MySQL error code, SQLSTATE value, and message string are available using C API functions:

- MySQL error code: Call `mysql_errno()`
- SQLSTATE value: Call `mysql_sqlstate()`
- Error message: Call `mysql_error()`

For prepared statements, the corresponding error functions are `mysql_stmt_errno()`, `mysql_stmt_sqlstate()`, and `mysql_stmt_error()`. All error functions are described in Section 27.7, “MySQL C API”.

The number of errors, warnings, and notes for the previous statement can be obtained by calling `mysql_warning_count()`. See Section 27.7.7.80, “mysql_warning_count()”.

The first two characters of an SQLSTATE value indicate the error class:

- Class = `'00'` indicates success.
- Class = `'01'` indicates a warning.
- Class = `'02'` indicates “not found.” This is relevant within the context of cursors and is used to control what happens when a cursor reaches the end of a data set. This condition also occurs for `SELECT ... INTO var_list` statements that retrieve no rows.
- Class > `'02'` indicates an exception.

B.3 Server Error Codes and Messages

MySQL programs have access to several types of error information when the server returns an error. For example, the `mysql` client program displays errors using the following format:

```
shell> SELECT * FROM no_such_table;
ERROR 1146 (42S02): Table 'test.no_such_table' doesn't exist
```

The message displayed contains three types of information:

- A numeric error code (`1146`). This number is MySQL-specific and is not portable to other database systems.
- A five-character SQLSTATE value (`'42S02'`). The values are taken from ANSI SQL and ODBC and are more standardized. Not all MySQL error numbers have corresponding SQLSTATE values. In these cases, `'HY000'` (general error) is used.
- A message string that provides a textual description of the error.

For error checking, use error codes, not error messages. Error messages do not change often, but it is possible. Also if the database administrator changes the language setting, that affects the language of error messages.

Error codes are stable across GA releases of a given MySQL series. Before a series reaches GA status, new codes may still be under development and subject to change.

Server error information comes from the following source files. For details about the way that error information is defined, see the MySQL Internals Manual.

- Error message information is listed in the `share/errmsg-utf8.txt` file. `%d` and `%s` represent numbers and strings, respectively, that are substituted into the Message values when they are displayed.
- The Error values listed in `share/errmsg-utf8.txt` are used to generate the definitions in the `include/mysqld_error.h` and `include/mysqld_ername.h` MySQL source files.
- The SQLSTATE values listed in `share/errmsg-utf8.txt` are used to generate the definitions in the `include/sql_state.h` MySQL source file.

Because updates are frequent, it is possible that those files will contain additional error information not listed here.

- Error: `1002` SQLSTATE: `HY000` (`ER_NO`)

  Message: NO

  Used in the construction of other messages.

- Error: `1003` SQLSTATE: `HY000` (`ER_YES`)

  Message: YES

  Used in the construction of other messages.

  Extended `EXPLAIN` format generates Note messages. `ER_YES` is used in the `Code` column for these messages in subsequent `SHOW WARNINGS` output.

- Error: `1004` SQLSTATE: `HY000` (`ER_CANT_CREATE_FILE`)

  Message: Can't create file '%s' (errno: %d - %s)

  Occurs for failure to create or copy a file needed for some operation.

  Possible causes: Permissions problem for source file; destination file already exists but is not writeable.

- Error: `1005` SQLSTATE: `HY000` (`ER_CANT_CREATE_TABLE`)

  Message: Can't create table '%s' (errno: %d - %s)

  `InnoDB` reports this error when a table cannot be created. If the error message refers to error 150, table creation failed because a foreign key constraint was not correctly formed. If the error message refers to error −1, table creation probably failed because the table includes a column name that matched the name of an internal `InnoDB` table.

- Error: `1006` SQLSTATE: `HY000` (`ER_CANT_CREATE_DB`)

  Message: Can't create database '%s' (errno: %d - %s)

- Error: `1007` SQLSTATE: `HY000` (`ER_DB_CREATE_EXISTS`)

  Message: Can't create database '%s'; database exists

  An attempt to create a database failed because the database already exists.

  Drop the database first if you really want to replace an existing database, or add an `IF NOT EXISTS` clause to the `CREATE DATABASE` statement if to retain an existing database without having the statement produce an error.

- Error: `1008` SQLSTATE: `HY000` (`ER_DB_DROP_EXISTS`)

  Message: Can't drop database '%s'; database doesn't exist

- Error: 1010 SQLSTATE: HY000 (ER_DB_DROP_RMDIR)

  Message: Error dropping database (can't rmdir '%s', errno: %d - %s)

- Error: 1012 SQLSTATE: HY000 (ER_CANT_FIND_SYSTEM_REC)

  Message: Can't read record in system table

  Returned by InnoDB for attempts to access InnoDB INFORMATION_SCHEMA tables when InnoDB is unavailable.

- Error: 1013 SQLSTATE: HY000 (ER_CANT_GET_STAT)

  Message: Can't get status of '%s' (errno: %d - %s)

- Error: 1015 SQLSTATE: HY000 (ER_CANT_LOCK)

  Message: Can't lock file (errno: %d - %s)

- Error: 1016 SQLSTATE: HY000 (ER_CANT_OPEN_FILE)

  Message: Can't open file: '%s' (errno: %d - %s)

  InnoDB reports this error when the table from the InnoDB data files cannot be found.

- Error: 1017 SQLSTATE: HY000 (ER_FILE_NOT_FOUND)

  Message: Can't find file: '%s' (errno: %d - %s)

- Error: 1018 SQLSTATE: HY000 (ER_CANT_READ_DIR)

  Message: Can't read dir of '%s' (errno: %d - %s)

- Error: 1020 SQLSTATE: HY000 (ER_CHECKREAD)

  Message: Record has changed since last read in table '%s'

- Error: 1022 SQLSTATE: 23000 (ER_DUP_KEY)

  Message: Can't write; duplicate key in table '%s'

- Error: 1024 SQLSTATE: HY000 (ER_ERROR_ON_READ)

  Message: Error reading file '%s' (errno: %d - %s)

- Error: 1025 SQLSTATE: HY000 (ER_ERROR_ON_RENAME)

  Message: Error on rename of '%s' to '%s' (errno: %d - %s)

- Error: 1026 SQLSTATE: HY000 (ER_ERROR_ON_WRITE)

  Message: Error writing file '%s' (errno: %d - %s)

- Error: 1027 SQLSTATE: HY000 (ER_FILE_USED)

  Message: '%s' is locked against change

- Error: 1028 SQLSTATE: HY000 (ER_FILSORT_ABORT)

  Message: Sort aborted

- Error: 1030 SQLSTATE: HY000 (ER_GET_ERRNO)

  Message: Got error %d - '%s' from storage engine

Check the %d value to see what the OS error means. For example, 28 indicates that you have run out of disk space.

- Error: 1031 SQLSTATE: HY000 (ER_ILLEGAL_HA)

  Message: Table storage engine for '%s' doesn't have this option

- Error: 1032 SQLSTATE: HY000 (ER_KEY_NOT_FOUND)

  Message: Can't find record in '%s'

- Error: 1033 SQLSTATE: HY000 (ER_NOT_FORM_FILE)

  Message: Incorrect information in file: '%s'

- Error: 1034 SQLSTATE: HY000 (ER_NOT_KEYFILE)

  Message: Incorrect key file for table '%s'; try to repair it

- Error: 1035 SQLSTATE: HY000 (ER_OLD_KEYFILE)

  Message: Old key file for table '%s'; repair it!

- Error: 1036 SQLSTATE: HY000 (ER_OPEN_AS_READONLY)

  Message: Table '%s' is read only

- Error: 1037 SQLSTATE: HY001 (ER_OUTOFMEMORY)

  Message: Out of memory; restart server and try again (needed %d bytes)

- Error: 1038 SQLSTATE: HY001 (ER_OUT_OF_SORTMEMORY)

  Message: Out of sort memory, consider increasing server sort buffer size

- Error: 1040 SQLSTATE: 08004 (ER_CON_COUNT_ERROR)

  Message: Too many connections

- Error: 1041 SQLSTATE: HY000 (ER_OUT_OF_RESOURCES)

  Message: Out of memory; check if mysqld or some other process uses all available memory; if not, you may have to use 'ulimit' to allow mysqld to use more memory or you can add more swap space

- Error: 1042 SQLSTATE: 08S01 (ER_BAD_HOST_ERROR)

  Message: Can't get hostname for your address

- Error: 1043 SQLSTATE: 08S01 (ER_HANDSHAKE_ERROR)

  Message: Bad handshake

- Error: 1044 SQLSTATE: 42000 (ER_DBACCESS_DENIED_ERROR)

  Message: Access denied for user '%s'@'%s' to database '%s'

- Error: 1045 SQLSTATE: 28000 (ER_ACCESS_DENIED_ERROR)

  Message: Access denied for user '%s'@'%s' (using password: %s)

- Error: 1046 SQLSTATE: 3D000 (ER_NO_DB_ERROR)

  Message: No database selected

- Error: 1047 SQLSTATE: 08S01 (ER_UNKNOWN_COM_ERROR)

Message: Unknown command

- Error: 1048 SQLSTATE: 23000 (ER_BAD_NULL_ERROR)

  Message: Column '%s' cannot be null

- Error: 1049 SQLSTATE: 42000 (ER_BAD_DB_ERROR)

  Message: Unknown database '%s'

- Error: 1050 SQLSTATE: 42S01 (ER_TABLE_EXISTS_ERROR)

  Message: Table '%s' already exists

- Error: 1051 SQLSTATE: 42S02 (ER_BAD_TABLE_ERROR)

  Message: Unknown table '%s'

- Error: 1052 SQLSTATE: 23000 (ER_NON_UNIQ_ERROR)

  Message: Column '%s' in %s is ambiguous

  ```
  %s = column name
  %s = location of column (for example, "field list")
  ```

  Likely cause: A column appears in a query without appropriate qualification, such as in a select list or ON clause.

  Examples:

  ```
  mysql> SELECT i FROM t INNER JOIN t AS t2;
  ERROR 1052 (23000): Column 'i' in field list is ambiguous

  mysql> SELECT * FROM t LEFT JOIN t AS t2 ON i = i;
  ERROR 1052 (23000): Column 'i' in on clause is ambiguous
  ```

  Resolution:

  - Qualify the column with the appropriate table name:

    ```
    mysql> SELECT t2.i FROM t INNER JOIN t AS t2;
    ```

  - Modify the query to avoid the need for qualification:

    ```
    mysql> SELECT * FROM t LEFT JOIN t AS t2 USING (i);
    ```

- Error: 1053 SQLSTATE: 08S01 (ER_SERVER_SHUTDOWN)

  Message: Server shutdown in progress

- Error: 1054 SQLSTATE: 42S22 (ER_BAD_FIELD_ERROR)

  Message: Unknown column '%s' in '%s'

- Error: 1055 SQLSTATE: 42000 (ER_WRONG_FIELD_WITH_GROUP)

  Message: '%s' isn't in GROUP BY

- Error: 1056 SQLSTATE: 42000 (ER_WRONG_GROUP_FIELD)

  Message: Can't group on '%s'

- Error: 1057 SQLSTATE: 42000 (ER_WRONG_SUM_SELECT)

  Message: Statement has sum functions and columns in same statement
- Error: 1058 SQLSTATE: 21S01 (ER_WRONG_VALUE_COUNT)

  Message: Column count doesn't match value count
- Error: 1059 SQLSTATE: 42000 (ER_TOO_LONG_IDENT)

  Message: Identifier name '%s' is too long
- Error: 1060 SQLSTATE: 42S21 (ER_DUP_FIELDNAME)

  Message: Duplicate column name '%s'
- Error: 1061 SQLSTATE: 42000 (ER_DUP_KEYNAME)

  Message: Duplicate key name '%s'
- Error: 1062 SQLSTATE: 23000 (ER_DUP_ENTRY)

  Message: Duplicate entry '%s' for key %d

  The message returned with this error uses the format string for ER_DUP_ENTRY_WITH_KEY_NAME.
- Error: 1063 SQLSTATE: 42000 (ER_WRONG_FIELD_SPEC)

  Message: Incorrect column specifier for column '%s'
- Error: 1064 SQLSTATE: 42000 (ER_PARSE_ERROR)

  Message: %s near '%s' at line %d
- Error: 1065 SQLSTATE: 42000 (ER_EMPTY_QUERY)

  Message: Query was empty
- Error: 1066 SQLSTATE: 42000 (ER_NONUNIQ_TABLE)

  Message: Not unique table/alias: '%s'
- Error: 1067 SQLSTATE: 42000 (ER_INVALID_DEFAULT)

  Message: Invalid default value for '%s'
- Error: 1068 SQLSTATE: 42000 (ER_MULTIPLE_PRI_KEY)

  Message: Multiple primary key defined
- Error: 1069 SQLSTATE: 42000 (ER_TOO_MANY_KEYS)

  Message: Too many keys specified; max %d keys allowed
- Error: 1070 SQLSTATE: 42000 (ER_TOO_MANY_KEY_PARTS)

  Message: Too many key parts specified; max %d parts allowed
- Error: 1071 SQLSTATE: 42000 (ER_TOO_LONG_KEY)

  Message: Specified key was too long; max key length is %d bytes
- Error: 1072 SQLSTATE: 42000 (ER_KEY_COLUMN_DOES_NOT_EXITS)

  Message: Key column '%s' doesn't exist in table

- Error: 1073 SQLSTATE: 42000 (ER_BLOB_USED_AS_KEY)

  Message: BLOB column '%s' can't be used in key specification with the used table type

- Error: 1074 SQLSTATE: 42000 (ER_TOO_BIG_FIELDLENGTH)

  Message: Column length too big for column '%s' (max = %lu); use BLOB or TEXT instead

- Error: 1075 SQLSTATE: 42000 (ER_WRONG_AUTO_KEY)

  Message: Incorrect table definition; there can be only one auto column and it must be defined as a key

- Error: 1076 SQLSTATE: HY000 (ER_READY)

  Message: %s: ready for connections. Version: '%s' socket: '%s' port: %d

- Error: 1077 SQLSTATE: HY000 (ER_NORMAL_SHUTDOWN)

  Message: %s: Normal shutdown

- Error: 1079 SQLSTATE: HY000 (ER_SHUTDOWN_COMPLETE)

  Message: %s: Shutdown complete

- Error: 1080 SQLSTATE: 08S01 (ER_FORCING_CLOSE)

  Message: %s: Forcing close of thread %ld user: '%s'

- Error: 1081 SQLSTATE: 08S01 (ER_IPSOCK_ERROR)

  Message: Can't create IP socket

- Error: 1082 SQLSTATE: 42S12 (ER_NO_SUCH_INDEX)

  Message: Table '%s' has no index like the one used in CREATE INDEX; recreate the table

- Error: 1083 SQLSTATE: 42000 (ER_WRONG_FIELD_TERMINATORS)

  Message: Field separator argument is not what is expected; check the manual

- Error: 1084 SQLSTATE: 42000 (ER_BLOBS_AND_NO_TERMINATED)

  Message: You can't use fixed rowlength with BLOBs; please use 'fields terminated by'

- Error: 1085 SQLSTATE: HY000 (ER_TEXTFILE_NOT_READABLE)

  Message: The file '%s' must be in the database directory or be readable by all

- Error: 1086 SQLSTATE: HY000 (ER_FILE_EXISTS_ERROR)

  Message: File '%s' already exists

- Error: 1087 SQLSTATE: HY000 (ER_LOAD_INFO)

  Message: Records: %ld Deleted: %ld Skipped: %ld Warnings: %ld

- Error: 1088 SQLSTATE: HY000 (ER_ALTER_INFO)

  Message: Records: %ld Duplicates: %ld

- Error: 1089 SQLSTATE: HY000 (ER_WRONG_SUB_KEY)

  Message: Incorrect prefix key; the used key part isn't a string, the used length is longer than the key part, or the storage engine doesn't support unique prefix keys

- Error: 1090 SQLSTATE: 42000 (ER_CANT_REMOVE_ALL_FIELDS)

  Message: You can't delete all columns with ALTER TABLE; use DROP TABLE instead

- Error: 1091 SQLSTATE: 42000 (ER_CANT_DROP_FIELD_OR_KEY)

  Message: Can't DROP '%s'; check that column/key exists

- Error: 1092 SQLSTATE: HY000 (ER_INSERT_INFO)

  Message: Records: %ld Duplicates: %ld Warnings: %ld

- Error: 1093 SQLSTATE: HY000 (ER_UPDATE_TABLE_USED)

  Message: You can't specify target table '%s' for update in FROM clause

  This error occurs for attempts to select from and modify the same table within a single statement. If the select attempt occurs within a derived table, you can avoid this error by setting the `derived_merge` flag of the `optimizer_switch` system variable to force the subquery to be materialized into a temporary table, which effectively causes it to be a different table from the one modified. See Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

- Error: 1094 SQLSTATE: HY000 (ER_NO_SUCH_THREAD)

  Message: Unknown thread id: %lu

- Error: 1095 SQLSTATE: HY000 (ER_KILL_DENIED_ERROR)

  Message: You are not owner of thread %lu

- Error: 1096 SQLSTATE: HY000 (ER_NO_TABLES_USED)

  Message: No tables used

- Error: 1097 SQLSTATE: HY000 (ER_TOO_BIG_SET)

  Message: Too many strings for column %s and SET

- Error: 1098 SQLSTATE: HY000 (ER_NO_UNIQUE_LOGFILE)

  Message: Can't generate a unique log-filename %s.(1-999)

- Error: 1099 SQLSTATE: HY000 (ER_TABLE_NOT_LOCKED_FOR_WRITE)

  Message: Table '%s' was locked with a READ lock and can't be updated

- Error: 1100 SQLSTATE: HY000 (ER_TABLE_NOT_LOCKED)

  Message: Table '%s' was not locked with LOCK TABLES

- Error: 1101 SQLSTATE: 42000 (ER_BLOB_CANT_HAVE_DEFAULT)

  Message: BLOB, TEXT, GEOMETRY or JSON column '%s' can't have a default value

- Error: 1102 SQLSTATE: 42000 (ER_WRONG_DB_NAME)

  Message: Incorrect database name '%s'

- Error: 1103 SQLSTATE: 42000 (ER_WRONG_TABLE_NAME)

  Message: Incorrect table name '%s'

- Error: 1104 SQLSTATE: 42000 (ER_TOO_BIG_SELECT)

Message: The SELECT would examine more than MAX_JOIN_SIZE rows; check your WHERE and use SET SQL_BIG_SELECTS=1 or SET MAX_JOIN_SIZE=# if the SELECT is okay

- Error: 1105 SQLSTATE: HY000 (ER_UNKNOWN_ERROR)

  Message: Unknown error

- Error: 1106 SQLSTATE: 42000 (ER_UNKNOWN_PROCEDURE)

  Message: Unknown procedure '%s'

- Error: 1107 SQLSTATE: 42000 (ER_WRONG_PARAMCOUNT_TO_PROCEDURE)

  Message: Incorrect parameter count to procedure '%s'

- Error: 1108 SQLSTATE: HY000 (ER_WRONG_PARAMETERS_TO_PROCEDURE)

  Message: Incorrect parameters to procedure '%s'

- Error: 1109 SQLSTATE: 42S02 (ER_UNKNOWN_TABLE)

  Message: Unknown table '%s' in %s

- Error: 1110 SQLSTATE: 42000 (ER_FIELD_SPECIFIED_TWICE)

  Message: Column '%s' specified twice

- Error: 1111 SQLSTATE: HY000 (ER_INVALID_GROUP_FUNC_USE)

  Message: Invalid use of group function

- Error: 1112 SQLSTATE: 42000 (ER_UNSUPPORTED_EXTENSION)

  Message: Table '%s' uses an extension that doesn't exist in this MySQL version

- Error: 1113 SQLSTATE: 42000 (ER_TABLE_MUST_HAVE_COLUMNS)

  Message: A table must have at least 1 column

- Error: 1114 SQLSTATE: HY000 (ER_RECORD_FILE_FULL)

  Message: The table '%s' is full

  InnoDB reports this error when the system tablespace runs out of free space. Reconfigure the system tablespace to add a new data file.

- Error: 1115 SQLSTATE: 42000 (ER_UNKNOWN_CHARACTER_SET)

  Message: Unknown character set: '%s'

- Error: 1116 SQLSTATE: HY000 (ER_TOO_MANY_TABLES)

  Message: Too many tables; MySQL can only use %d tables in a join

- Error: 1117 SQLSTATE: HY000 (ER_TOO_MANY_FIELDS)

  Message: Too many columns

- Error: 1118 SQLSTATE: 42000 (ER_TOO_BIG_ROWSIZE)

  Message: Row size too large. The maximum row size for the used table type, not counting BLOBs, is %ld. This includes storage overhead, check the manual. You have to change some columns to TEXT or BLOBs

- Error: 1119 SQLSTATE: HY000 (ER_STACK_OVERRUN)

  Message: Thread stack overrun: Used: %ld of a %ld stack. Use 'mysqld --thread_stack=#' to specify a bigger stack if needed

- Error: 1120 SQLSTATE: 42000 (ER_WRONG_OUTER_JOIN)

  Message: Cross dependency found in OUTER JOIN; examine your ON conditions

  ER_WRONG_OUTER_JOIN was removed after 8.0.0.

- Error: 1120 SQLSTATE: 42000 (ER_WRONG_OUTER_JOIN_UNUSED)

  Message: Cross dependency found in OUTER JOIN; examine your ON conditions

  ER_WRONG_OUTER_JOIN_UNUSED was added in 8.0.1.

- Error: 1121 SQLSTATE: 42000 (ER_NULL_COLUMN_IN_INDEX)

  Message: Table handler doesn't support NULL in given index. Please change column '%s' to be NOT NULL or use another handler

- Error: 1122 SQLSTATE: HY000 (ER_CANT_FIND_UDF)

  Message: Can't load function '%s'

- Error: 1123 SQLSTATE: HY000 (ER_CANT_INITIALIZE_UDF)

  Message: Can't initialize function '%s'; %s

- Error: 1124 SQLSTATE: HY000 (ER_UDF_NO_PATHS)

  Message: No paths allowed for shared library

- Error: 1125 SQLSTATE: HY000 (ER_UDF_EXISTS)

  Message: Function '%s' already exists

- Error: 1126 SQLSTATE: HY000 (ER_CANT_OPEN_LIBRARY)

  Message: Can't open shared library '%s' (errno: %d %s)

- Error: 1127 SQLSTATE: HY000 (ER_CANT_FIND_DL_ENTRY)

  Message: Can't find symbol '%s' in library

- Error: 1128 SQLSTATE: HY000 (ER_FUNCTION_NOT_DEFINED)

  Message: Function '%s' is not defined

- Error: 1129 SQLSTATE: HY000 (ER_HOST_IS_BLOCKED)

  Message: Host '%s' is blocked because of many connection errors; unblock with 'mysqladmin flush-hosts'

- Error: 1130 SQLSTATE: HY000 (ER_HOST_NOT_PRIVILEGED)

  Message: Host '%s' is not allowed to connect to this MySQL server

- Error: 1131 SQLSTATE: 42000 (ER_PASSWORD_ANONYMOUS_USER)

  Message: You are using MySQL as an anonymous user and anonymous users are not allowed to change passwords

- Error: 1132 SQLSTATE: 42000 (ER_PASSWORD_NOT_ALLOWED)

  Message: You must have privileges to update tables in the mysql database to be able to change passwords for others

- Error: 1133 SQLSTATE: 42000 (ER_PASSWORD_NO_MATCH)

  Message: Can't find any matching row in the user table

- Error: 1134 SQLSTATE: HY000 (ER_UPDATE_INFO)

  Message: Rows matched: %ld Changed: %ld Warnings: %ld

- Error: 1135 SQLSTATE: HY000 (ER_CANT_CREATE_THREAD)

  Message: Can't create a new thread (errno %d); if you are not out of available memory, you can consult the manual for a possible OS-dependent bug

- Error: 1136 SQLSTATE: 21S01 (ER_WRONG_VALUE_COUNT_ON_ROW)

  Message: Column count doesn't match value count at row %ld

- Error: 1137 SQLSTATE: HY000 (ER_CANT_REOPEN_TABLE)

  Message: Can't reopen table: '%s'

- Error: 1138 SQLSTATE: 22004 (ER_INVALID_USE_OF_NULL)

  Message: Invalid use of NULL value

- Error: 1139 SQLSTATE: 42000 (ER_REGEXP_ERROR)

  Message: Got error '%s' from regexp

- Error: 1140 SQLSTATE: 42000 (ER_MIX_OF_GROUP_FUNC_AND_FIELDS)

  Message: Mixing of GROUP columns (MIN(),MAX(),COUNT(),...) with no GROUP columns is illegal if there is no GROUP BY clause

- Error: 1141 SQLSTATE: 42000 (ER_NONEXISTING_GRANT)

  Message: There is no such grant defined for user '%s' on host '%s'

- Error: 1142 SQLSTATE: 42000 (ER_TABLEACCESS_DENIED_ERROR)

  Message: %s command denied to user '%s'@'%s' for table '%s'

- Error: 1143 SQLSTATE: 42000 (ER_COLUMNACCESS_DENIED_ERROR)

  Message: %s command denied to user '%s'@'%s' for column '%s' in table '%s'

- Error: 1144 SQLSTATE: 42000 (ER_ILLEGAL_GRANT_FOR_TABLE)

  Message: Illegal GRANT/REVOKE command; please consult the manual to see which privileges can be used

- Error: 1145 SQLSTATE: 42000 (ER_GRANT_WRONG_HOST_OR_USER)

  Message: The host or user argument to GRANT is too long

- Error: 1146 SQLSTATE: 42S02 (ER_NO_SUCH_TABLE)

  Message: Table '%s.%s' doesn't exist

- Error: 1147 SQLSTATE: 42000 (ER_NONEXISTING_TABLE_GRANT)

  Message: There is no such grant defined for user '%s' on host '%s' on table '%s'

- Error: 1148 SQLSTATE: 42000 (ER_NOT_ALLOWED_COMMAND)

  Message: The used command is not allowed with this MySQL version

- Error: 1149 SQLSTATE: 42000 (ER_SYNTAX_ERROR)

  Message: You have an error in your SQL syntax; check the manual that corresponds to your MySQL server version for the right syntax to use

- Error: 1152 SQLSTATE: 08S01 (ER_ABORTING_CONNECTION)

  Message: Aborted connection %ld to db: '%s' user: '%s' (%s)

- Error: 1153 SQLSTATE: 08S01 (ER_NET_PACKET_TOO_LARGE)

  Message: Got a packet bigger than 'max_allowed_packet' bytes

- Error: 1154 SQLSTATE: 08S01 (ER_NET_READ_ERROR_FROM_PIPE)

  Message: Got a read error from the connection pipe

- Error: 1155 SQLSTATE: 08S01 (ER_NET_FCNTL_ERROR)

  Message: Got an error from fcntl()

- Error: 1156 SQLSTATE: 08S01 (ER_NET_PACKETS_OUT_OF_ORDER)

  Message: Got packets out of order

- Error: 1157 SQLSTATE: 08S01 (ER_NET_UNCOMPRESS_ERROR)

  Message: Couldn't uncompress communication packet

- Error: 1158 SQLSTATE: 08S01 (ER_NET_READ_ERROR)

  Message: Got an error reading communication packets

- Error: 1159 SQLSTATE: 08S01 (ER_NET_READ_INTERRUPTED)

  Message: Got timeout reading communication packets

- Error: 1160 SQLSTATE: 08S01 (ER_NET_ERROR_ON_WRITE)

  Message: Got an error writing communication packets

- Error: 1161 SQLSTATE: 08S01 (ER_NET_WRITE_INTERRUPTED)

  Message: Got timeout writing communication packets

- Error: 1162 SQLSTATE: 42000 (ER_TOO_LONG_STRING)

  Message: Result string is longer than 'max_allowed_packet' bytes

- Error: 1163 SQLSTATE: 42000 (ER_TABLE_CANT_HANDLE_BLOB)

  Message: The used table type doesn't support BLOB/TEXT columns

- Error: 1164 SQLSTATE: 42000 (ER_TABLE_CANT_HANDLE_AUTO_INCREMENT)

  Message: The used table type doesn't support AUTO_INCREMENT columns

- Error: 1166 SQLSTATE: 42000 (ER_WRONG_COLUMN_NAME)

  Message: Incorrect column name '%s'

- Error: 1167 SQLSTATE: 42000 (ER_WRONG_KEY_COLUMN)

  Message: The used storage engine can't index column '%s'

- Error: 1168 SQLSTATE: HY000 (ER_WRONG_MRG_TABLE)

  Message: Unable to open underlying table which is differently defined or of non-MyISAM type or doesn't exist

- Error: 1169 SQLSTATE: 23000 (ER_DUP_UNIQUE)

  Message: Can't write, because of unique constraint, to table '%s'

- Error: 1170 SQLSTATE: 42000 (ER_BLOB_KEY_WITHOUT_LENGTH)

  Message: BLOB/TEXT column '%s' used in key specification without a key length

- Error: 1171 SQLSTATE: 42000 (ER_PRIMARY_CANT_HAVE_NULL)

  Message: All parts of a PRIMARY KEY must be NOT NULL; if you need NULL in a key, use UNIQUE instead

- Error: 1172 SQLSTATE: 42000 (ER_TOO_MANY_ROWS)

  Message: Result consisted of more than one row

- Error: 1173 SQLSTATE: 42000 (ER_REQUIRES_PRIMARY_KEY)

  Message: This table type requires a primary key

- Error: 1175 SQLSTATE: HY000 (ER_UPDATE_WITHOUT_KEY_IN_SAFE_MODE)

  Message: You are using safe update mode and you tried to update a table without a WHERE that uses a KEY column

- Error: 1176 SQLSTATE: 42000 (ER_KEY_DOES_NOT_EXITS)

  Message: Key '%s' doesn't exist in table '%s'

- Error: 1177 SQLSTATE: 42000 (ER_CHECK_NO_SUCH_TABLE)

  Message: Can't open table

- Error: 1178 SQLSTATE: 42000 (ER_CHECK_NOT_IMPLEMENTED)

  Message: The storage engine for the table doesn't support %s

- Error: 1179 SQLSTATE: 25000 (ER_CANT_DO_THIS_DURING_AN_TRANSACTION)

  Message: You are not allowed to execute this command in a transaction

- Error: 1180 SQLSTATE: HY000 (ER_ERROR_DURING_COMMIT)

  Message: Got error %d - '%s' during COMMIT

- Error: 1181 SQLSTATE: HY000 (ER_ERROR_DURING_ROLLBACK)

  Message: Got error %d - '%s' during ROLLBACK

- Error: 1182 SQLSTATE: HY000 (ER_ERROR_DURING_FLUSH_LOGS)

Message: Got error %d during FLUSH_LOGS

- Error: 1184 SQLSTATE: 08S01 (ER_NEW_ABORTING_CONNECTION)

  Message: Aborted connection %u to db: '%s' user: '%s' host: '%s' (%s)

- Error: 1188 SQLSTATE: HY000 (ER_MASTER)

  Message: Error from master: '%s'

- Error: 1189 SQLSTATE: 08S01 (ER_MASTER_NET_READ)

  Message: Net error reading from master

- Error: 1190 SQLSTATE: 08S01 (ER_MASTER_NET_WRITE)

  Message: Net error writing to master

- Error: 1191 SQLSTATE: HY000 (ER_FT_MATCHING_KEY_NOT_FOUND)

  Message: Can't find FULLTEXT index matching the column list

- Error: 1192 SQLSTATE: HY000 (ER_LOCK_OR_ACTIVE_TRANSACTION)

  Message: Can't execute the given command because you have active locked tables or an active transaction

- Error: 1193 SQLSTATE: HY000 (ER_UNKNOWN_SYSTEM_VARIABLE)

  Message: Unknown system variable '%s'

- Error: 1194 SQLSTATE: HY000 (ER_CRASHED_ON_USAGE)

  Message: Table '%s' is marked as crashed and should be repaired

- Error: 1195 SQLSTATE: HY000 (ER_CRASHED_ON_REPAIR)

  Message: Table '%s' is marked as crashed and last (automatic?) repair failed

- Error: 1196 SQLSTATE: HY000 (ER_WARNING_NOT_COMPLETE_ROLLBACK)

  Message: Some non-transactional changed tables couldn't be rolled back

- Error: 1197 SQLSTATE: HY000 (ER_TRANS_CACHE_FULL)

  Message: Multi-statement transaction required more than 'max_binlog_cache_size' bytes of storage; increase this mysqld variable and try again

- Error: 1199 SQLSTATE: HY000 (ER_SLAVE_NOT_RUNNING)

  Message: This operation requires a running slave; configure slave and do START SLAVE

- Error: 1200 SQLSTATE: HY000 (ER_BAD_SLAVE)

  Message: The server is not configured as slave; fix in config file or with CHANGE MASTER TO

- Error: 1201 SQLSTATE: HY000 (ER_MASTER_INFO)

  Message: Could not initialize master info structure; more error messages can be found in the MySQL error log

- Error: 1202 SQLSTATE: HY000 (ER_SLAVE_THREAD)

  Message: Could not create slave thread; check system resources

- Error: 1203 SQLSTATE: 42000 (ER_TOO_MANY_USER_CONNECTIONS)

  Message: User %s already has more than 'max_user_connections' active connections

- Error: 1204 SQLSTATE: HY000 (ER_SET_CONSTANTS_ONLY)

  Message: You may only use constant expressions with SET

- Error: 1205 SQLSTATE: HY000 (ER_LOCK_WAIT_TIMEOUT)

  Message: Lock wait timeout exceeded; try restarting transaction

  InnoDB reports this error when lock wait timeout expires. The statement that waited too long was rolled back (not the entire transaction). You can increase the value of the innodb_lock_wait_timeout configuration option if SQL statements should wait longer for other transactions to complete, or decrease it if too many long-running transactions are causing locking problems and reducing concurrency on a busy system.

- Error: 1206 SQLSTATE: HY000 (ER_LOCK_TABLE_FULL)

  Message: The total number of locks exceeds the lock table size

  InnoDB reports this error when the total number of locks exceeds the amount of memory devoted to managing locks. To avoid this error, increase the value of innodb_buffer_pool_size. Within an individual application, a workaround may be to break a large operation into smaller pieces. For example, if the error occurs for a large INSERT, perform several smaller INSERT operations.

- Error: 1207 SQLSTATE: 25000 (ER_READ_ONLY_TRANSACTION)

  Message: Update locks cannot be acquired during a READ UNCOMMITTED transaction

- Error: 1210 SQLSTATE: HY000 (ER_WRONG_ARGUMENTS)

  Message: Incorrect arguments to %s

- Error: 1211 SQLSTATE: 42000 (ER_NO_PERMISSION_TO_CREATE_USER)

  Message: '%s'@'%s' is not allowed to create new users

- Error: 1213 SQLSTATE: 40001 (ER_LOCK_DEADLOCK)

  Message: Deadlock found when trying to get lock; try restarting transaction

  InnoDB reports this error when a transaction encounters a deadlock and is automatically rolled back so that your application can take corrective action. To recover from this error, run all the operations in this transaction again. A deadlock occurs when requests for locks arrive in inconsistent order between transactions. The transaction that was rolled back released all its locks, and the other transaction can now get all the locks it requested. Thus, when you re-run the transaction that was rolled back, it might have to wait for other transactions to complete, but typically the deadlock does not recur. If you encounter frequent deadlocks, make the sequence of locking operations (LOCK TABLES, SELECT ... FOR UPDATE, and so on) consistent between the different transactions or applications that experience the issue. See Section 15.5.5, “Deadlocks in InnoDB” for details.

- Error: 1214 SQLSTATE: HY000 (ER_TABLE_CANT_HANDLE_FT)

  Message: The used table type doesn't support FULLTEXT indexes

- Error: 1215 SQLSTATE: HY000 (ER_CANNOT_ADD_FOREIGN)

  Message: Cannot add foreign key constraint

- Error: 1216 SQLSTATE: 23000 (ER_NO_REFERENCED_ROW)

Message: Cannot add or update a child row: a foreign key constraint fails

`InnoDB` reports this error when you try to add a row but there is no parent row, and a foreign key constraint fails. Add the parent row first.

- Error: `1217` SQLSTATE: `23000` (`ER_ROW_IS_REFERENCED`)

  Message: Cannot delete or update a parent row: a foreign key constraint fails

  `InnoDB` reports this error when you try to delete a parent row that has children, and a foreign key constraint fails. Delete the children first.

- Error: `1218` SQLSTATE: `08S01` (`ER_CONNECT_TO_MASTER`)

  Message: Error connecting to master: %s

- Error: `1220` SQLSTATE: `HY000` (`ER_ERROR_WHEN_EXECUTING_COMMAND`)

  Message: Error when executing command %s: %s

- Error: `1221` SQLSTATE: `HY000` (`ER_WRONG_USAGE`)

  Message: Incorrect usage of %s and %s

- Error: `1222` SQLSTATE: `21000` (`ER_WRONG_NUMBER_OF_COLUMNS_IN_SELECT`)

  Message: The used SELECT statements have a different number of columns

- Error: `1223` SQLSTATE: `HY000` (`ER_CANT_UPDATE_WITH_READLOCK`)

  Message: Can't execute the query because you have a conflicting read lock

- Error: `1224` SQLSTATE: `HY000` (`ER_MIXING_NOT_ALLOWED`)

  Message: Mixing of transactional and non-transactional tables is disabled

- Error: `1225` SQLSTATE: `HY000` (`ER_DUP_ARGUMENT`)

  Message: Option '%s' used twice in statement

- Error: `1226` SQLSTATE: `42000` (`ER_USER_LIMIT_REACHED`)

  Message: User '%s' has exceeded the '%s' resource (current value: %ld)

- Error: `1227` SQLSTATE: `42000` (`ER_SPECIFIC_ACCESS_DENIED_ERROR`)

  Message: Access denied; you need (at least one of) the %s privilege(s) for this operation

- Error: `1228` SQLSTATE: `HY000` (`ER_LOCAL_VARIABLE`)

  Message: Variable '%s' is a SESSION variable and can't be used with SET GLOBAL

- Error: `1229` SQLSTATE: `HY000` (`ER_GLOBAL_VARIABLE`)

  Message: Variable '%s' is a GLOBAL variable and should be set with SET GLOBAL

- Error: `1230` SQLSTATE: `42000` (`ER_NO_DEFAULT`)

  Message: Variable '%s' doesn't have a default value

- Error: `1231` SQLSTATE: `42000` (`ER_WRONG_VALUE_FOR_VAR`)

  Message: Variable '%s' can't be set to the value of '%s'

- Error: 1232 SQLSTATE: 42000 (ER_WRONG_TYPE_FOR_VAR)

  Message: Incorrect argument type to variable '%s'
- Error: 1233 SQLSTATE: HY000 (ER_VAR_CANT_BE_READ)

  Message: Variable '%s' can only be set, not read
- Error: 1234 SQLSTATE: 42000 (ER_CANT_USE_OPTION_HERE)

  Message: Incorrect usage/placement of '%s'
- Error: 1235 SQLSTATE: 42000 (ER_NOT_SUPPORTED_YET)

  Message: This version of MySQL doesn't yet support '%s'
- Error: 1236 SQLSTATE: HY000 (ER_MASTER_FATAL_ERROR_READING_BINLOG)

  Message: Got fatal error %d from master when reading data from binary log: '%s'
- Error: 1237 SQLSTATE: HY000 (ER_SLAVE_IGNORED_TABLE)

  Message: Slave SQL thread ignored the query because of replicate-*-table rules
- Error: 1238 SQLSTATE: HY000 (ER_INCORRECT_GLOBAL_LOCAL_VAR)

  Message: Variable '%s' is a %s variable
- Error: 1239 SQLSTATE: 42000 (ER_WRONG_FK_DEF)

  Message: Incorrect foreign key definition for '%s': %s
- Error: 1240 SQLSTATE: HY000 (ER_KEY_REF_DO_NOT_MATCH_TABLE_REF)

  Message: Key reference and table reference don't match
- Error: 1241 SQLSTATE: 21000 (ER_OPERAND_COLUMNS)

  Message: Operand should contain %d column(s)
- Error: 1242 SQLSTATE: 21000 (ER_SUBQUERY_NO_1_ROW)

  Message: Subquery returns more than 1 row
- Error: 1243 SQLSTATE: HY000 (ER_UNKNOWN_STMT_HANDLER)

  Message: Unknown prepared statement handler (%.*s) given to %s
- Error: 1244 SQLSTATE: HY000 (ER_CORRUPT_HELP_DB)

  Message: Help database is corrupt or does not exist
- Error: 1246 SQLSTATE: HY000 (ER_AUTO_CONVERT)

  Message: Converting column '%s' from %s to %s
- Error: 1247 SQLSTATE: 42S22 (ER_ILLEGAL_REFERENCE)

  Message: Reference '%s' not supported (%s)
- Error: 1248 SQLSTATE: 42000 (ER_DERIVED_MUST_HAVE_ALIAS)

  Message: Every derived table must have its own alias
- Error: 1249 SQLSTATE: 01000 (ER_SELECT_REDUCED)

Message: Select %u was reduced during optimization

- Error: 1250 SQLSTATE: 42000 (ER_TABLENAME_NOT_ALLOWED_HERE)

  Message: Table '%s' from one of the SELECTs cannot be used in %s

- Error: 1251 SQLSTATE: 08004 (ER_NOT_SUPPORTED_AUTH_MODE)

  Message: Client does not support authentication protocol requested by server; consider upgrading MySQL client

- Error: 1252 SQLSTATE: 42000 (ER_SPATIAL_CANT_HAVE_NULL)

  Message: All parts of a SPATIAL index must be NOT NULL

- Error: 1253 SQLSTATE: 42000 (ER_COLLATION_CHARSET_MISMATCH)

  Message: COLLATION '%s' is not valid for CHARACTER SET '%s'

- Error: 1256 SQLSTATE: HY000 (ER_TOO_BIG_FOR_UNCOMPRESS)

  Message: Uncompressed data size too large; the maximum size is %d (probably, length of uncompressed data was corrupted)

- Error: 1257 SQLSTATE: HY000 (ER_ZLIB_Z_MEM_ERROR)

  Message: ZLIB: Not enough memory

- Error: 1258 SQLSTATE: HY000 (ER_ZLIB_Z_BUF_ERROR)

  Message: ZLIB: Not enough room in the output buffer (probably, length of uncompressed data was corrupted)

- Error: 1259 SQLSTATE: HY000 (ER_ZLIB_Z_DATA_ERROR)

  Message: ZLIB: Input data corrupted

- Error: 1260 SQLSTATE: HY000 (ER_CUT_VALUE_GROUP_CONCAT)

  Message: Row %u was cut by GROUP_CONCAT()

- Error: 1261 SQLSTATE: 01000 (ER_WARN_TOO_FEW_RECORDS)

  Message: Row %ld doesn't contain data for all columns

- Error: 1262 SQLSTATE: 01000 (ER_WARN_TOO_MANY_RECORDS)

  Message: Row %ld was truncated; it contained more data than there were input columns

- Error: 1263 SQLSTATE: 22004 (ER_WARN_NULL_TO_NOTNULL)

  Message: Column set to default value; NULL supplied to NOT NULL column '%s' at row %ld

- Error: 1264 SQLSTATE: 22003 (ER_WARN_DATA_OUT_OF_RANGE)

  Message: Out of range value for column '%s' at row %ld

- Error: 1265 SQLSTATE: 01000 (WARN_DATA_TRUNCATED)

  Message: Data truncated for column '%s' at row %ld

- Error: 1266 SQLSTATE: HY000 (ER_WARN_USING_OTHER_HANDLER)

  Message: Using storage engine %s for table '%s'

- Error: 1267 SQLSTATE: HY000 (ER_CANT_AGGREGATE_2COLLATIONS)

  Message: Illegal mix of collations (%s,%s) and (%s,%s) for operation '%s'

- Error: 1269 SQLSTATE: HY000 (ER_REVOKE_GRANTS)

  Message: Can't revoke all privileges for one or more of the requested users

- Error: 1270 SQLSTATE: HY000 (ER_CANT_AGGREGATE_3COLLATIONS)

  Message: Illegal mix of collations (%s,%s), (%s,%s), (%s,%s) for operation '%s'

- Error: 1271 SQLSTATE: HY000 (ER_CANT_AGGREGATE_NCOLLATIONS)

  Message: Illegal mix of collations for operation '%s'

- Error: 1272 SQLSTATE: HY000 (ER_VARIABLE_IS_NOT_STRUCT)

  Message: Variable '%s' is not a variable component (can't be used as XXXX.variable_name)

- Error: 1273 SQLSTATE: HY000 (ER_UNKNOWN_COLLATION)

  Message: Unknown collation: '%s'

- Error: 1274 SQLSTATE: HY000 (ER_SLAVE_IGNORED_SSL_PARAMS)

  Message: SSL parameters in CHANGE MASTER are ignored because this MySQL slave was compiled without SSL support; they can be used later if MySQL slave with SSL is started

- Error: 1275 SQLSTATE: HY000 (ER_SERVER_IS_IN_SECURE_AUTH_MODE)

  Message: Server is running in --secure-auth mode, but '%s'@'%s' has a password in the old format; please change the password to the new format

- Error: 1276 SQLSTATE: HY000 (ER_WARN_FIELD_RESOLVED)

  Message: Field or reference '%s%s%s%s%s' of SELECT #%d was resolved in SELECT #%d

- Error: 1277 SQLSTATE: HY000 (ER_BAD_SLAVE_UNTIL_COND)

  Message: Incorrect parameter or combination of parameters for START SLAVE UNTIL

- Error: 1278 SQLSTATE: HY000 (ER_MISSING_SKIP_SLAVE)

  Message: It is recommended to use --skip-slave-start when doing step-by-step replication with START SLAVE UNTIL; otherwise, you will get problems if you get an unexpected slave's mysqld restart

- Error: 1279 SQLSTATE: HY000 (ER_UNTIL_COND_IGNORED)

  Message: SQL thread is not to be started so UNTIL options are ignored

- Error: 1280 SQLSTATE: 42000 (ER_WRONG_NAME_FOR_INDEX)

  Message: Incorrect index name '%s'

- Error: 1281 SQLSTATE: 42000 (ER_WRONG_NAME_FOR_CATALOG)

  Message: Incorrect catalog name '%s'

- Error: 1282 SQLSTATE: HY000 (ER_WARN_QC_RESIZE)

  Message: Query cache failed to set size %lu; new query cache size is %lu

- Error: 1283 SQLSTATE: HY000 (ER_BAD_FT_COLUMN)

  Message: Column '%s' cannot be part of FULLTEXT index

- Error: 1284 SQLSTATE: HY000 (ER_UNKNOWN_KEY_CACHE)

  Message: Unknown key cache '%s'

- Error: 1285 SQLSTATE: HY000 (ER_WARN_HOSTNAME_WONT_WORK)

  Message: MySQL is started in --skip-name-resolve mode; you must restart it without this switch for this grant to work

- Error: 1286 SQLSTATE: 42000 (ER_UNKNOWN_STORAGE_ENGINE)

  Message: Unknown storage engine '%s'

- Error: 1287 SQLSTATE: HY000 (ER_WARN_DEPRECATED_SYNTAX)

  Message: '%s' is deprecated and will be removed in a future release. Please use %s instead

- Error: 1288 SQLSTATE: HY000 (ER_NON_UPDATABLE_TABLE)

  Message: The target table %s of the %s is not updatable

- Error: 1289 SQLSTATE: HY000 (ER_FEATURE_DISABLED)

  Message: The '%s' feature is disabled; you need MySQL built with '%s' to have it working

- Error: 1290 SQLSTATE: HY000 (ER_OPTION_PREVENTS_STATEMENT)

  Message: The MySQL server is running with the %s option so it cannot execute this statement

- Error: 1291 SQLSTATE: HY000 (ER_DUPLICATED_VALUE_IN_TYPE)

  Message: Column '%s' has duplicated value '%s' in %s

- Error: 1292 SQLSTATE: 22007 (ER_TRUNCATED_WRONG_VALUE)

  Message: Truncated incorrect %s value: '%s'

- Error: 1294 SQLSTATE: HY000 (ER_INVALID_ON_UPDATE)

  Message: Invalid ON UPDATE clause for '%s' column

- Error: 1295 SQLSTATE: HY000 (ER_UNSUPPORTED_PS)

  Message: This command is not supported in the prepared statement protocol yet

- Error: 1296 SQLSTATE: HY000 (ER_GET_ERRMSG)

  Message: Got error %d '%s' from %s

- Error: 1297 SQLSTATE: HY000 (ER_GET_TEMPORARY_ERRMSG)

  Message: Got temporary error %d '%s' from %s

- Error: 1298 SQLSTATE: HY000 (ER_UNKNOWN_TIME_ZONE)

  Message: Unknown or incorrect time zone: '%s'

- Error: 1299 SQLSTATE: HY000 (ER_WARN_INVALID_TIMESTAMP)

  Message: Invalid TIMESTAMP value in column '%s' at row %ld

- Error: 1300 SQLSTATE: HY000 (ER_INVALID_CHARACTER_STRING)

  Message: Invalid %s character string: '%s'

- Error: 1301 SQLSTATE: HY000 (ER_WARN_ALLOWED_PACKET_OVERFLOWED)

  Message: Result of %s() was larger than max_allowed_packet (%ld) - truncated

- Error: 1302 SQLSTATE: HY000 (ER_CONFLICTING_DECLARATIONS)

  Message: Conflicting declarations: '%s%s' and '%s%s'

- Error: 1303 SQLSTATE: 2F003 (ER_SP_NO_RECURSIVE_CREATE)

  Message: Can't create a %s from within another stored routine

- Error: 1304 SQLSTATE: 42000 (ER_SP_ALREADY_EXISTS)

  Message: %s %s already exists

- Error: 1305 SQLSTATE: 42000 (ER_SP_DOES_NOT_EXIST)

  Message: %s %s does not exist

- Error: 1306 SQLSTATE: HY000 (ER_SP_DROP_FAILED)

  Message: Failed to DROP %s %s

- Error: 1307 SQLSTATE: HY000 (ER_SP_STORE_FAILED)

  Message: Failed to CREATE %s %s

- Error: 1308 SQLSTATE: 42000 (ER_SP_LILABEL_MISMATCH)

  Message: %s with no matching label: %s

- Error: 1309 SQLSTATE: 42000 (ER_SP_LABEL_REDEFINE)

  Message: Redefining label %s

- Error: 1310 SQLSTATE: 42000 (ER_SP_LABEL_MISMATCH)

  Message: End-label %s without match

- Error: 1311 SQLSTATE: 01000 (ER_SP_UNINIT_VAR)

  Message: Referring to uninitialized variable %s

- Error: 1312 SQLSTATE: 0A000 (ER_SP_BADSELECT)

  Message: PROCEDURE %s can't return a result set in the given context

- Error: 1313 SQLSTATE: 42000 (ER_SP_BADRETURN)

  Message: RETURN is only allowed in a FUNCTION

- Error: 1314 SQLSTATE: 0A000 (ER_SP_BADSTATEMENT)

  Message: %s is not allowed in stored procedures

- Error: 1315 SQLSTATE: 42000 (ER_UPDATE_LOG_DEPRECATED_IGNORED)

  Message: The update log is deprecated and replaced by the binary log; SET SQL_LOG_UPDATE has been ignored.

- Error: 1316 SQLSTATE: 42000 (ER_UPDATE_LOG_DEPRECATED_TRANSLATED)

  Message: The update log is deprecated and replaced by the binary log; SET SQL_LOG_UPDATE has been translated to SET SQL_LOG_BIN.

- Error: 1317 SQLSTATE: 70100 (ER_QUERY_INTERRUPTED)

  Message: Query execution was interrupted

- Error: 1318 SQLSTATE: 42000 (ER_SP_WRONG_NO_OF_ARGS)

  Message: Incorrect number of arguments for %s %s; expected %u, got %u

- Error: 1319 SQLSTATE: 42000 (ER_SP_COND_MISMATCH)

  Message: Undefined CONDITION: %s

- Error: 1320 SQLSTATE: 42000 (ER_SP_NORETURN)

  Message: No RETURN found in FUNCTION %s

- Error: 1321 SQLSTATE: 2F005 (ER_SP_NORETURNEND)

  Message: FUNCTION %s ended without RETURN

- Error: 1322 SQLSTATE: 42000 (ER_SP_BAD_CURSOR_QUERY)

  Message: Cursor statement must be a SELECT

- Error: 1323 SQLSTATE: 42000 (ER_SP_BAD_CURSOR_SELECT)

  Message: Cursor SELECT must not have INTO

- Error: 1324 SQLSTATE: 42000 (ER_SP_CURSOR_MISMATCH)

  Message: Undefined CURSOR: %s

- Error: 1325 SQLSTATE: 24000 (ER_SP_CURSOR_ALREADY_OPEN)

  Message: Cursor is already open

- Error: 1326 SQLSTATE: 24000 (ER_SP_CURSOR_NOT_OPEN)

  Message: Cursor is not open

- Error: 1327 SQLSTATE: 42000 (ER_SP_UNDECLARED_VAR)

  Message: Undeclared variable: %s

- Error: 1328 SQLSTATE: HY000 (ER_SP_WRONG_NO_OF_FETCH_ARGS)

  Message: Incorrect number of FETCH variables

- Error: 1329 SQLSTATE: 02000 (ER_SP_FETCH_NO_DATA)

  Message: No data - zero rows fetched, selected, or processed

- Error: 1330 SQLSTATE: 42000 (ER_SP_DUP_PARAM)

  Message: Duplicate parameter: %s

- Error: 1331 SQLSTATE: 42000 (ER_SP_DUP_VAR)

  Message: Duplicate variable: %s

- Error: 1332 SQLSTATE: 42000 (ER_SP_DUP_COND)

  Message: Duplicate condition: %s

- Error: 1333 SQLSTATE: 42000 (ER_SP_DUP_CURS)

  Message: Duplicate cursor: %s

- Error: 1334 SQLSTATE: HY000 (ER_SP_CANT_ALTER)

  Message: Failed to ALTER %s %s

- Error: 1335 SQLSTATE: 0A000 (ER_SP_SUBSELECT_NYI)

  Message: Subquery value not supported

- Error: 1336 SQLSTATE: 0A000 (ER_STMT_NOT_ALLOWED_IN_SF_OR_TRG)

  Message: %s is not allowed in stored function or trigger

- Error: 1337 SQLSTATE: 42000 (ER_SP_VARCOND_AFTER_CURSHNDLR)

  Message: Variable or condition declaration after cursor or handler declaration

- Error: 1338 SQLSTATE: 42000 (ER_SP_CURSOR_AFTER_HANDLER)

  Message: Cursor declaration after handler declaration

- Error: 1339 SQLSTATE: 20000 (ER_SP_CASE_NOT_FOUND)

  Message: Case not found for CASE statement

- Error: 1340 SQLSTATE: HY000 (ER_FPARSER_TOO_BIG_FILE)

  Message: Configuration file '%s' is too big

- Error: 1341 SQLSTATE: HY000 (ER_FPARSER_BAD_HEADER)

  Message: Malformed file type header in file '%s'

- Error: 1342 SQLSTATE: HY000 (ER_FPARSER_EOF_IN_COMMENT)

  Message: Unexpected end of file while parsing comment '%s'

- Error: 1343 SQLSTATE: HY000 (ER_FPARSER_ERROR_IN_PARAMETER)

  Message: Error while parsing parameter '%s' (line: '%s')

- Error: 1344 SQLSTATE: HY000 (ER_FPARSER_EOF_IN_UNKNOWN_PARAMETER)

  Message: Unexpected end of file while skipping unknown parameter '%s'

- Error: 1345 SQLSTATE: HY000 (ER_VIEW_NO_EXPLAIN)

  Message: EXPLAIN/SHOW can not be issued; lacking privileges for underlying table

- Error: 1347 SQLSTATE: HY000 (ER_WRONG_OBJECT)

  Message: '%s.%s' is not %s

  The named object is incorrect for the type of operation attempted on it. It must be an object of the named type. Example: `HANDLER OPEN` requires a base table, not a view. It fails if attempted on an `INFORMATION_SCHEMA` table that is implemented as a view on data dictionary tables.

- Error: 1348 SQLSTATE: HY000 (ER_NONUPDATEABLE_COLUMN)

Message: Column '%s' is not updatable

- Error: 1350 SQLSTATE: HY000 (ER_VIEW_SELECT_CLAUSE)

  Message: View's SELECT contains a '%s' clause

- Error: 1351 SQLSTATE: HY000 (ER_VIEW_SELECT_VARIABLE)

  Message: View's SELECT contains a variable or parameter

- Error: 1352 SQLSTATE: HY000 (ER_VIEW_SELECT_TMPTABLE)

  Message: View's SELECT refers to a temporary table '%s'

- Error: 1353 SQLSTATE: HY000 (ER_VIEW_WRONG_LIST)

  Message: View's SELECT and view's field list have different column counts

- Error: 1354 SQLSTATE: HY000 (ER_WARN_VIEW_MERGE)

  Message: View merge algorithm can't be used here for now (assumed undefined algorithm)

- Error: 1355 SQLSTATE: HY000 (ER_WARN_VIEW_WITHOUT_KEY)

  Message: View being updated does not have complete key of underlying table in it

- Error: 1356 SQLSTATE: HY000 (ER_VIEW_INVALID)

  Message: View '%s.%s' references invalid table(s) or column(s) or function(s) or definer/invoker of view lack rights to use them

- Error: 1357 SQLSTATE: HY000 (ER_SP_NO_DROP_SP)

  Message: Can't drop or alter a %s from within another stored routine

- Error: 1359 SQLSTATE: HY000 (ER_TRG_ALREADY_EXISTS)

  Message: Trigger already exists

- Error: 1360 SQLSTATE: HY000 (ER_TRG_DOES_NOT_EXIST)

  Message: Trigger does not exist

- Error: 1361 SQLSTATE: HY000 (ER_TRG_ON_VIEW_OR_TEMP_TABLE)

  Message: Trigger's '%s' is view or temporary table

- Error: 1362 SQLSTATE: HY000 (ER_TRG_CANT_CHANGE_ROW)

  Message: Updating of %s row is not allowed in %strigger

- Error: 1363 SQLSTATE: HY000 (ER_TRG_NO_SUCH_ROW_IN_TRG)

  Message: There is no %s row in %s trigger

- Error: 1364 SQLSTATE: HY000 (ER_NO_DEFAULT_FOR_FIELD)

  Message: Field '%s' doesn't have a default value

- Error: 1365 SQLSTATE: 22012 (ER_DIVISION_BY_ZERO)

  Message: Division by 0

- Error: 1366 SQLSTATE: HY000 (ER_TRUNCATED_WRONG_VALUE_FOR_FIELD)

Message: Incorrect %s value: '%s' for column '%s' at row %ld

- Error: 1367 SQLSTATE: 22007 (ER_ILLEGAL_VALUE_FOR_TYPE)

  Message: Illegal %s '%s' value found during parsing

- Error: 1368 SQLSTATE: HY000 (ER_VIEW_NONUPD_CHECK)

  Message: CHECK OPTION on non-updatable view '%s.%s'

- Error: 1369 SQLSTATE: HY000 (ER_VIEW_CHECK_FAILED)

  Message: CHECK OPTION failed '%s.%s'

- Error: 1370 SQLSTATE: 42000 (ER_PROCACCESS_DENIED_ERROR)

  Message: %s command denied to user '%s'@'%s' for routine '%s'

- Error: 1371 SQLSTATE: HY000 (ER_RELAY_LOG_FAIL)

  Message: Failed purging old relay logs: %s

- Error: 1373 SQLSTATE: HY000 (ER_UNKNOWN_TARGET_BINLOG)

  Message: Target log not found in binlog index

- Error: 1374 SQLSTATE: HY000 (ER_IO_ERR_LOG_INDEX_READ)

  Message: I/O error reading log index file

- Error: 1375 SQLSTATE: HY000 (ER_BINLOG_PURGE_PROHIBITED)

  Message: Server configuration does not permit binlog purge

- Error: 1376 SQLSTATE: HY000 (ER_FSEEK_FAIL)

  Message: Failed on fseek()

- Error: 1377 SQLSTATE: HY000 (ER_BINLOG_PURGE_FATAL_ERR)

  Message: Fatal error during log purge

- Error: 1378 SQLSTATE: HY000 (ER_LOG_IN_USE)

  Message: A purgeable log is in use, will not purge

- Error: 1379 SQLSTATE: HY000 (ER_LOG_PURGE_UNKNOWN_ERR)

  Message: Unknown error during log purge

- Error: 1380 SQLSTATE: HY000 (ER_RELAY_LOG_INIT)

  Message: Failed initializing relay log position: %s

- Error: 1381 SQLSTATE: HY000 (ER_NO_BINARY_LOGGING)

  Message: You are not using binary logging

- Error: 1382 SQLSTATE: HY000 (ER_RESERVED_SYNTAX)

  Message: The '%s' syntax is reserved for purposes internal to the MySQL server

- Error: 1390 SQLSTATE: HY000 (ER_PS_MANY_PARAM)

  Message: Prepared statement contains too many placeholders

- Error: 1391 SQLSTATE: HY000 (ER_KEY_PART_0)

  Message: Key part '%s' length cannot be 0

- Error: 1392 SQLSTATE: HY000 (ER_VIEW_CHECKSUM)

  Message: View text checksum failed

- Error: 1393 SQLSTATE: HY000 (ER_VIEW_MULTIUPDATE)

  Message: Can not modify more than one base table through a join view '%s.%s'

- Error: 1394 SQLSTATE: HY000 (ER_VIEW_NO_INSERT_FIELD_LIST)

  Message: Can not insert into join view '%s.%s' without fields list

- Error: 1395 SQLSTATE: HY000 (ER_VIEW_DELETE_MERGE_VIEW)

  Message: Can not delete from join view '%s.%s'

- Error: 1396 SQLSTATE: HY000 (ER_CANNOT_USER)

  Message: Operation %s failed for %s

- Error: 1397 SQLSTATE: XAE04 (ER_XAER_NOTA)

  Message: XAER_NOTA: Unknown XID

- Error: 1398 SQLSTATE: XAE05 (ER_XAER_INVAL)

  Message: XAER_INVAL: Invalid arguments (or unsupported command)

- Error: 1399 SQLSTATE: XAE07 (ER_XAER_RMFAIL)

  Message: XAER_RMFAIL: The command cannot be executed when global transaction is in the %s state

- Error: 1400 SQLSTATE: XAE09 (ER_XAER_OUTSIDE)

  Message: XAER_OUTSIDE: Some work is done outside global transaction

- Error: 1401 SQLSTATE: XAE03 (ER_XAER_RMERR)

  Message: XAER_RMERR: Fatal error occurred in the transaction branch - check your data for consistency

- Error: 1402 SQLSTATE: XA100 (ER_XA_RBROLLBACK)

  Message: XA_RBROLLBACK: Transaction branch was rolled back

- Error: 1403 SQLSTATE: 42000 (ER_NONEXISTING_PROC_GRANT)

  Message: There is no such grant defined for user '%s' on host '%s' on routine '%s'

- Error: 1404 SQLSTATE: HY000 (ER_PROC_AUTO_GRANT_FAIL)

  Message: Failed to grant EXECUTE and ALTER ROUTINE privileges

- Error: 1405 SQLSTATE: HY000 (ER_PROC_AUTO_REVOKE_FAIL)

  Message: Failed to revoke all privileges to dropped routine

- Error: 1406 SQLSTATE: 22001 (ER_DATA_TOO_LONG)

  Message: Data too long for column '%s' at row %ld

- Error: 1407 SQLSTATE: 42000 (ER_SP_BAD_SQLSTATE)

  Message: Bad SQLSTATE: '%s'

- Error: 1408 SQLSTATE: HY000 (ER_STARTUP)

  Message: %s: ready for connections. Version: '%s' socket: '%s' port: %d %s

- Error: 1409 SQLSTATE: HY000 (ER_LOAD_FROM_FIXED_SIZE_ROWS_TO_VAR)

  Message: Can't load value from file with fixed size rows to variable

- Error: 1410 SQLSTATE: 42000 (ER_CANT_CREATE_USER_WITH_GRANT)

  Message: You are not allowed to create a user with GRANT

- Error: 1411 SQLSTATE: HY000 (ER_WRONG_VALUE_FOR_TYPE)

  Message: Incorrect %s value: '%s' for function %s

- Error: 1412 SQLSTATE: HY000 (ER_TABLE_DEF_CHANGED)

  Message: Table definition has changed, please retry transaction

- Error: 1413 SQLSTATE: 42000 (ER_SP_DUP_HANDLER)

  Message: Duplicate handler declared in the same block

- Error: 1414 SQLSTATE: 42000 (ER_SP_NOT_VAR_ARG)

  Message: OUT or INOUT argument %d for routine %s is not a variable or NEW pseudo-variable in BEFORE trigger

- Error: 1415 SQLSTATE: 0A000 (ER_SP_NO_RETSET)

  Message: Not allowed to return a result set from a %s

- Error: 1416 SQLSTATE: 22003 (ER_CANT_CREATE_GEOMETRY_OBJECT)

  Message: Cannot get geometry object from data you send to the GEOMETRY field

- Error: 1418 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_ROUTINE)

  Message: This function has none of DETERMINISTIC, NO SQL, or READS SQL DATA in its declaration and binary logging is enabled (you *might* want to use the less safe log_bin_trust_function_creators variable)

- Error: 1419 SQLSTATE: HY000 (ER_BINLOG_CREATE_ROUTINE_NEED_SUPER)

  Message: You do not have the SUPER privilege and binary logging is enabled (you *might* want to use the less safe log_bin_trust_function_creators variable)

- Error: 1421 SQLSTATE: HY000 (ER_STMT_HAS_NO_OPEN_CURSOR)

  Message: The statement (%lu) has no open cursor.

- Error: 1422 SQLSTATE: HY000 (ER_COMMIT_NOT_ALLOWED_IN_SF_OR_TRG)

  Message: Explicit or implicit commit is not allowed in stored function or trigger.

- Error: 1423 SQLSTATE: HY000 (ER_NO_DEFAULT_FOR_VIEW_FIELD)

  Message: Field of view '%s.%s' underlying table doesn't have a default value

- Error: 1424 SQLSTATE: HY000 (ER_SP_NO_RECURSION)

Message: Recursive stored functions and triggers are not allowed.

- Error: 1425 SQLSTATE: 42000 (ER_TOO_BIG_SCALE)

  Message: Too big scale %d specified for column '%s'. Maximum is %lu.

- Error: 1426 SQLSTATE: 42000 (ER_TOO_BIG_PRECISION)

  Message: Too-big precision %d specified for '%s'. Maximum is %lu.

- Error: 1427 SQLSTATE: 42000 (ER_M_BIGGER_THAN_D)

  Message: For float(M,D), double(M,D) or decimal(M,D), M must be >= D (column '%s').

- Error: 1428 SQLSTATE: HY000 (ER_WRONG_LOCK_OF_SYSTEM_TABLE)

  Message: You can't combine write-locking of system tables with other tables or lock types

- Error: 1429 SQLSTATE: HY000 (ER_CONNECT_TO_FOREIGN_DATA_SOURCE)

  Message: Unable to connect to foreign data source: %s

- Error: 1430 SQLSTATE: HY000 (ER_QUERY_ON_FOREIGN_DATA_SOURCE)

  Message: There was a problem processing the query on the foreign data source. Data source error: %s

- Error: 1431 SQLSTATE: HY000 (ER_FOREIGN_DATA_SOURCE_DOESNT_EXIST)

  Message: The foreign data source you are trying to reference does not exist. Data source error: %s

- Error: 1432 SQLSTATE: HY000 (ER_FOREIGN_DATA_STRING_INVALID_CANT_CREATE)

  Message: Can't create federated table. The data source connection string '%s' is not in the correct format

- Error: 1433 SQLSTATE: HY000 (ER_FOREIGN_DATA_STRING_INVALID)

  Message: The data source connection string '%s' is not in the correct format

- Error: 1435 SQLSTATE: HY000 (ER_TRG_IN_WRONG_SCHEMA)

  Message: Trigger in wrong schema

- Error: 1436 SQLSTATE: HY000 (ER_STACK_OVERRUN_NEED_MORE)

  Message: Thread stack overrun: %ld bytes used of a %ld byte stack, and %ld bytes needed. Use 'mysqld --thread_stack=#' to specify a bigger stack.

- Error: 1437 SQLSTATE: 42000 (ER_TOO_LONG_BODY)

  Message: Routine body for '%s' is too long

- Error: 1438 SQLSTATE: HY000 (ER_WARN_CANT_DROP_DEFAULT_KEYCACHE)

  Message: Cannot drop default keycache

- Error: 1439 SQLSTATE: 42000 (ER_TOO_BIG_DISPLAYWIDTH)

  Message: Display width out of range for column '%s' (max = %lu)

- Error: 1440 SQLSTATE: XAE08 (ER_XAER_DUPID)

  Message: XAER_DUPID: The XID already exists

- Error: 1441 SQLSTATE: 22008 (ER_DATETIME_FUNCTION_OVERFLOW)

  Message: Datetime function: %s field overflow

- Error: 1442 SQLSTATE: HY000 (ER_CANT_UPDATE_USED_TABLE_IN_SF_OR_TRG)

  Message: Can't update table '%s' in stored function/trigger because it is already used by statement which invoked this stored function/trigger.

- Error: 1443 SQLSTATE: HY000 (ER_VIEW_PREVENT_UPDATE)

  Message: The definition of table '%s' prevents operation %s on table '%s'.

- Error: 1444 SQLSTATE: HY000 (ER_PS_NO_RECURSION)

  Message: The prepared statement contains a stored routine call that refers to that same statement. It's not allowed to execute a prepared statement in such a recursive manner

- Error: 1445 SQLSTATE: HY000 (ER_SP_CANT_SET_AUTOCOMMIT)

  Message: Not allowed to set autocommit from a stored function or trigger

- Error: 1447 SQLSTATE: HY000 (ER_VIEW_FRM_NO_USER)

  Message: View '%s'.'%s' has no definer information (old table format). Current user is used as definer. Please recreate the view!

- Error: 1448 SQLSTATE: HY000 (ER_VIEW_OTHER_USER)

  Message: You need the SUPER privilege for creation view with '%s'@'%s' definer

- Error: 1449 SQLSTATE: HY000 (ER_NO_SUCH_USER)

  Message: The user specified as a definer ('%s'@'%s') does not exist

- Error: 1450 SQLSTATE: HY000 (ER_FORBID_SCHEMA_CHANGE)

  Message: Changing schema from '%s' to '%s' is not allowed.

- Error: 1451 SQLSTATE: 23000 (ER_ROW_IS_REFERENCED_2)

  Message: Cannot delete or update a parent row: a foreign key constraint fails (%s)

- Error: 1452 SQLSTATE: 23000 (ER_NO_REFERENCED_ROW_2)

  Message: Cannot add or update a child row: a foreign key constraint fails (%s)

- Error: 1453 SQLSTATE: 42000 (ER_SP_BAD_VAR_SHADOW)

  Message: Variable '%s' must be quoted with `...`, or renamed

- Error: 1454 SQLSTATE: HY000 (ER_TRG_NO_DEFINER)

  Message: No definer attribute for trigger '%s'.'%s'. It's disallowed to create trigger without definer.

- Error: 1455 SQLSTATE: HY000 (ER_OLD_FILE_FORMAT)

  Message: '%s' has an old format, you should re-create the '%s' object(s)

- Error: 1456 SQLSTATE: HY000 (ER_SP_RECURSION_LIMIT)

  Message: Recursive limit %d (as set by the max_sp_recursion_depth variable) was exceeded for routine %s

- Error: 1458 SQLSTATE: 42000 (ER_SP_WRONG_NAME)

  Message: Incorrect routine name '%s'

- Error: 1459 SQLSTATE: HY000 (ER_TABLE_NEEDS_UPGRADE)

  Message: Table upgrade required. Please do "REPAIR TABLE `%s`" or dump/reload to fix it!

- Error: 1460 SQLSTATE: 42000 (ER_SP_NO_AGGREGATE)

  Message: AGGREGATE is not supported for stored functions

- Error: 1461 SQLSTATE: 42000 (ER_MAX_PREPARED_STMT_COUNT_REACHED)

  Message: Can't create more than max_prepared_stmt_count statements (current value: %lu)

- Error: 1462 SQLSTATE: HY000 (ER_VIEW_RECURSIVE)

  Message: `%s`.`%s` contains view recursion

- Error: 1463 SQLSTATE: 42000 (ER_NON_GROUPING_FIELD_USED)

  Message: Non-grouping field '%s' is used in %s clause

- Error: 1464 SQLSTATE: HY000 (ER_TABLE_CANT_HANDLE_SPKEYS)

  Message: The used table type doesn't support SPATIAL indexes

- Error: 1465 SQLSTATE: HY000 (ER_NO_TRIGGERS_ON_SYSTEM_SCHEMA)

  Message: Triggers can not be created on system tables

- Error: 1466 SQLSTATE: HY000 (ER_REMOVED_SPACES)

  Message: Leading spaces are removed from name '%s'

- Error: 1467 SQLSTATE: HY000 (ER_AUTOINC_READ_FAILED)

  Message: Failed to read auto-increment value from storage engine

- Error: 1468 SQLSTATE: HY000 (ER_USERNAME)

  Message: user name

- Error: 1469 SQLSTATE: HY000 (ER_HOSTNAME)

  Message: host name

- Error: 1470 SQLSTATE: HY000 (ER_WRONG_STRING_LENGTH)

  Message: String '%s' is too long for %s (should be no longer than %d)

- Error: 1471 SQLSTATE: HY000 (ER_NON_INSERTABLE_TABLE)

  Message: The target table %s of the %s is not insertable-into

- Error: 1472 SQLSTATE: HY000 (ER_ADMIN_WRONG_MRG_TABLE)

  Message: Table '%s' is differently defined or of non-MyISAM type or doesn't exist

- Error: 1473 SQLSTATE: HY000 (ER_TOO_HIGH_LEVEL_OF_NESTING_FOR_SELECT)

  Message: Too high level of nesting for select

- Error: 1474 SQLSTATE: HY000 (ER_NAME_BECOMES_EMPTY)

Message: Name '%s' has become ''

- Error: 1475 SQLSTATE: HY000 (ER_AMBIGUOUS_FIELD_TERM)

  Message: First character of the FIELDS TERMINATED string is ambiguous; please use non-optional and non-empty FIELDS ENCLOSED BY

- Error: 1476 SQLSTATE: HY000 (ER_FOREIGN_SERVER_EXISTS)

  Message: The foreign server, %s, you are trying to create already exists.

- Error: 1477 SQLSTATE: HY000 (ER_FOREIGN_SERVER_DOESNT_EXIST)

  Message: The foreign server name you are trying to reference does not exist. Data source error: %s

- Error: 1478 SQLSTATE: HY000 (ER_ILLEGAL_HA_CREATE_OPTION)

  Message: Table storage engine '%s' does not support the create option '%s'

- Error: 1479 SQLSTATE: HY000 (ER_PARTITION_REQUIRES_VALUES_ERROR)

  Message: Syntax error: %s PARTITIONING requires definition of VALUES %s for each partition

- Error: 1480 SQLSTATE: HY000 (ER_PARTITION_WRONG_VALUES_ERROR)

  Message: Only %s PARTITIONING can use VALUES %s in partition definition

- Error: 1481 SQLSTATE: HY000 (ER_PARTITION_MAXVALUE_ERROR)

  Message: MAXVALUE can only be used in last partition definition

- Error: 1484 SQLSTATE: HY000 (ER_PARTITION_WRONG_NO_PART_ERROR)

  Message: Wrong number of partitions defined, mismatch with previous setting

- Error: 1485 SQLSTATE: HY000 (ER_PARTITION_WRONG_NO_SUBPART_ERROR)

  Message: Wrong number of subpartitions defined, mismatch with previous setting

- Error: 1486 SQLSTATE: HY000 (ER_WRONG_EXPR_IN_PARTITION_FUNC_ERROR)

  Message: Constant, random or timezone-dependent expressions in (sub)partitioning function are not allowed

- Error: 1488 SQLSTATE: HY000 (ER_FIELD_NOT_FOUND_PART_ERROR)

  Message: Field in list of fields for partition function not found in table

- Error: 1490 SQLSTATE: HY000 (ER_INCONSISTENT_PARTITION_INFO_ERROR)

  Message: The partition info in the frm file is not consistent with what can be written into the frm file

- Error: 1491 SQLSTATE: HY000 (ER_PARTITION_FUNC_NOT_ALLOWED_ERROR)

  Message: The %s function returns the wrong type

- Error: 1492 SQLSTATE: HY000 (ER_PARTITIONS_MUST_BE_DEFINED_ERROR)

  Message: For %s partitions each partition must be defined

- Error: 1493 SQLSTATE: HY000 (ER_RANGE_NOT_INCREASING_ERROR)

  Message: VALUES LESS THAN value must be strictly increasing for each partition

- Error: 1494 SQLSTATE: HY000 (ER_INCONSISTENT_TYPE_OF_FUNCTIONS_ERROR)

  Message: VALUES value must be of same type as partition function

- Error: 1495 SQLSTATE: HY000 (ER_MULTIPLE_DEF_CONST_IN_LIST_PART_ERROR)

  Message: Multiple definition of same constant in list partitioning

- Error: 1496 SQLSTATE: HY000 (ER_PARTITION_ENTRY_ERROR)

  Message: Partitioning can not be used stand-alone in query

- Error: 1497 SQLSTATE: HY000 (ER_MIX_HANDLER_ERROR)

  Message: The mix of handlers in the partitions is not allowed in this version of MySQL

- Error: 1498 SQLSTATE: HY000 (ER_PARTITION_NOT_DEFINED_ERROR)

  Message: For the partitioned engine it is necessary to define all %s

- Error: 1499 SQLSTATE: HY000 (ER_TOO_MANY_PARTITIONS_ERROR)

  Message: Too many partitions (including subpartitions) were defined

- Error: 1500 SQLSTATE: HY000 (ER_SUBPARTITION_ERROR)

  Message: It is only possible to mix RANGE/LIST partitioning with HASH/KEY partitioning for subpartitioning

- Error: 1501 SQLSTATE: HY000 (ER_CANT_CREATE_HANDLER_FILE)

  Message: Failed to create specific handler file

- Error: 1502 SQLSTATE: HY000 (ER_BLOB_FIELD_IN_PART_FUNC_ERROR)

  Message: A BLOB field is not allowed in partition function

- Error: 1503 SQLSTATE: HY000 (ER_UNIQUE_KEY_NEED_ALL_FIELDS_IN_PF)

  Message: A %s must include all columns in the table's partitioning function

- Error: 1504 SQLSTATE: HY000 (ER_NO_PARTS_ERROR)

  Message: Number of %s = 0 is not an allowed value

- Error: 1505 SQLSTATE: HY000 (ER_PARTITION_MGMT_ON_NONPARTITIONED)

  Message: Partition management on a not partitioned table is not possible

- Error: 1506 SQLSTATE: HY000 (ER_FOREIGN_KEY_ON_PARTITIONED)

  Message: Foreign keys are not yet supported in conjunction with partitioning

- Error: 1507 SQLSTATE: HY000 (ER_DROP_PARTITION_NON_EXISTENT)

  Message: Error in list of partitions to %s

- Error: 1508 SQLSTATE: HY000 (ER_DROP_LAST_PARTITION)

  Message: Cannot remove all partitions, use DROP TABLE instead

- Error: 1509 SQLSTATE: HY000 (ER_COALESCE_ONLY_ON_HASH_PARTITION)

  Message: COALESCE PARTITION can only be used on HASH/KEY partitions

- Error: 1510 SQLSTATE: HY000 (ER_REORG_HASH_ONLY_ON_SAME_NO)

  Message: REORGANIZE PARTITION can only be used to reorganize partitions not to change their numbers

- Error: 1511 SQLSTATE: HY000 (ER_REORG_NO_PARAM_ERROR)

  Message: REORGANIZE PARTITION without parameters can only be used on auto-partitioned tables using HASH PARTITIONs

- Error: 1512 SQLSTATE: HY000 (ER_ONLY_ON_RANGE_LIST_PARTITION)

  Message: %s PARTITION can only be used on RANGE/LIST partitions

- Error: 1513 SQLSTATE: HY000 (ER_ADD_PARTITION_SUBPART_ERROR)

  Message: Trying to Add partition(s) with wrong number of subpartitions

- Error: 1514 SQLSTATE: HY000 (ER_ADD_PARTITION_NO_NEW_PARTITION)

  Message: At least one partition must be added

- Error: 1515 SQLSTATE: HY000 (ER_COALESCE_PARTITION_NO_PARTITION)

  Message: At least one partition must be coalesced

- Error: 1516 SQLSTATE: HY000 (ER_REORG_PARTITION_NOT_EXIST)

  Message: More partitions to reorganize than there are partitions

- Error: 1517 SQLSTATE: HY000 (ER_SAME_NAME_PARTITION)

  Message: Duplicate partition name %s

- Error: 1518 SQLSTATE: HY000 (ER_NO_BINLOG_ERROR)

  Message: It is not allowed to shut off binlog on this command

- Error: 1519 SQLSTATE: HY000 (ER_CONSECUTIVE_REORG_PARTITIONS)

  Message: When reorganizing a set of partitions they must be in consecutive order

- Error: 1520 SQLSTATE: HY000 (ER_REORG_OUTSIDE_RANGE)

  Message: Reorganize of range partitions cannot change total ranges except for last partition where it can extend the range

- Error: 1521 SQLSTATE: HY000 (ER_PARTITION_FUNCTION_FAILURE)

  Message: Partition function not supported in this version for this handler

- Error: 1523 SQLSTATE: HY000 (ER_LIMITED_PART_RANGE)

  Message: The %s handler only supports 32 bit integers in VALUES

- Error: 1524 SQLSTATE: HY000 (ER_PLUGIN_IS_NOT_LOADED)

  Message: Plugin '%s' is not loaded

- Error: 1525 SQLSTATE: HY000 (ER_WRONG_VALUE)

  Message: Incorrect %s value: '%s'

- Error: 1526 SQLSTATE: HY000 (ER_NO_PARTITION_FOR_GIVEN_VALUE)

Message: Table has no partition for value %s

- Error: 1527 SQLSTATE: HY000 (ER_FILEGROUP_OPTION_ONLY_ONCE)

  Message: It is not allowed to specify %s more than once

- Error: 1528 SQLSTATE: HY000 (ER_CREATE_FILEGROUP_FAILED)

  Message: Failed to create %s

- Error: 1529 SQLSTATE: HY000 (ER_DROP_FILEGROUP_FAILED)

  Message: Failed to drop %s

- Error: 1530 SQLSTATE: HY000 (ER_TABLESPACE_AUTO_EXTEND_ERROR)

  Message: The handler doesn't support autoextend of tablespaces

- Error: 1531 SQLSTATE: HY000 (ER_WRONG_SIZE_NUMBER)

  Message: A size parameter was incorrectly specified, either number or on the form 10M

- Error: 1532 SQLSTATE: HY000 (ER_SIZE_OVERFLOW_ERROR)

  Message: The size number was correct but we don't allow the digit part to be more than 2 billion

- Error: 1533 SQLSTATE: HY000 (ER_ALTER_FILEGROUP_FAILED)

  Message: Failed to alter: %s

- Error: 1534 SQLSTATE: HY000 (ER_BINLOG_ROW_LOGGING_FAILED)

  Message: Writing one row to the row-based binary log failed

- Error: 1537 SQLSTATE: HY000 (ER_EVENT_ALREADY_EXISTS)

  Message: Event '%s' already exists

- Error: 1539 SQLSTATE: HY000 (ER_EVENT_DOES_NOT_EXIST)

  Message: Unknown event '%s'

- Error: 1542 SQLSTATE: HY000 (ER_EVENT_INTERVAL_NOT_POSITIVE_OR_TOO_BIG)

  Message: INTERVAL is either not positive or too big

- Error: 1543 SQLSTATE: HY000 (ER_EVENT_ENDS_BEFORE_STARTS)

  Message: ENDS is either invalid or before STARTS

- Error: 1544 SQLSTATE: HY000 (ER_EVENT_EXEC_TIME_IN_THE_PAST)

  Message: Event execution time is in the past. Event has been disabled

- Error: 1551 SQLSTATE: HY000 (ER_EVENT_SAME_NAME)

  Message: Same old and new event name

- Error: 1553 SQLSTATE: HY000 (ER_DROP_INDEX_FK)

  Message: Cannot drop index '%s': needed in a foreign key constraint

  InnoDB reports this error when you attempt to drop the last index that can enforce a particular referential constraint.

For optimal performance with DML statements, InnoDB requires an index to exist on foreign key columns, so that UPDATE and DELETE operations on a parent table can easily check whether corresponding rows exist in the child table. MySQL creates or drops such indexes automatically when needed, as a side-effect of CREATE TABLE, CREATE INDEX, and ALTER TABLE statements.

When you drop an index, InnoDB checks if the index is used for checking a foreign key constraint. It is still OK to drop the index if there is another index that can be used to enforce the same constraint. InnoDB prevents you from dropping the last index that can enforce a particular referential constraint.

- Error: 1554 SQLSTATE: HY000 (ER_WARN_DEPRECATED_SYNTAX_WITH_VER)

  Message: The syntax '%s' is deprecated and will be removed in MySQL %s. Please use %s instead

- Error: 1556 SQLSTATE: HY000 (ER_CANT_LOCK_LOG_TABLE)

  Message: You can't use locks with log tables.

- Error: 1557 SQLSTATE: 23000 (ER_FOREIGN_DUPLICATE_KEY_OLD_UNUSED)

  Message: Upholding foreign key constraints for table '%s', entry '%s', key %d would lead to a duplicate entry

- Error: 1558 SQLSTATE: HY000 (ER_COL_COUNT_DOESNT_MATCH_PLEASE_UPDATE)

  Message: Column count of mysql.%s is wrong. Expected %d, found %d. Created with MySQL %d, now running %d. Please use mysql_upgrade to fix this error.

- Error: 1559 SQLSTATE: HY000 (ER_TEMP_TABLE_PREVENTS_SWITCH_OUT_OF_RBR)

  Message: Cannot switch out of the row-based binary log format when the session has open temporary tables

- Error: 1560 SQLSTATE: HY000 (ER_STORED_FUNCTION_PREVENTS_SWITCH_BINLOG_FORMAT)

  Message: Cannot change the binary logging format inside a stored function or trigger

- Error: 1562 SQLSTATE: HY000 (ER_PARTITION_NO_TEMPORARY)

  Message: Cannot create temporary table with partitions

- Error: 1563 SQLSTATE: HY000 (ER_PARTITION_CONST_DOMAIN_ERROR)

  Message: Partition constant is out of partition function domain

- Error: 1564 SQLSTATE: HY000 (ER_PARTITION_FUNCTION_IS_NOT_ALLOWED)

  Message: This partition function is not allowed

- Error: 1565 SQLSTATE: HY000 (ER_DDL_LOG_ERROR)

  Message: Error in DDL log

  ER_DDL_LOG_ERROR was removed after 8.0.1.

- Error: 1566 SQLSTATE: HY000 (ER_NULL_IN_VALUES_LESS_THAN)

  Message: Not allowed to use NULL value in VALUES LESS THAN

- Error: 1567 SQLSTATE: HY000 (ER_WRONG_PARTITION_NAME)

  Message: Incorrect partition name

- Error: 1568 SQLSTATE: 25001 (ER_CANT_CHANGE_TX_CHARACTERISTICS)

Message: Transaction characteristics can't be changed while a transaction is in progress

- Error: 1569 SQLSTATE: HY000 (ER_DUP_ENTRY_AUTOINCREMENT_CASE)

  Message: ALTER TABLE causes auto_increment resequencing, resulting in duplicate entry '%s' for key '%s'

- Error: 1571 SQLSTATE: HY000 (ER_EVENT_SET_VAR_ERROR)

  Message: Error during starting/stopping of the scheduler. Error code %u

- Error: 1572 SQLSTATE: HY000 (ER_PARTITION_MERGE_ERROR)

  Message: Engine cannot be used in partitioned tables

- Error: 1575 SQLSTATE: HY000 (ER_BASE64_DECODE_ERROR)

  Message: Decoding of base64 string failed

- Error: 1576 SQLSTATE: HY000 (ER_EVENT_RECURSION_FORBIDDEN)

  Message: Recursion of EVENT DDL statements is forbidden when body is present

- Error: 1578 SQLSTATE: HY000 (ER_ONLY_INTEGERS_ALLOWED)

  Message: Only integers allowed as number here

- Error: 1579 SQLSTATE: HY000 (ER_UNSUPORTED_LOG_ENGINE)

  Message: This storage engine cannot be used for log tables

- Error: 1580 SQLSTATE: HY000 (ER_BAD_LOG_STATEMENT)

  Message: You cannot '%s' a log table if logging is enabled

- Error: 1581 SQLSTATE: HY000 (ER_CANT_RENAME_LOG_TABLE)

  Message: Cannot rename '%s'. When logging enabled, rename to/from log table must rename two tables: the log table to an archive table and another table back to '%s'

- Error: 1582 SQLSTATE: 42000 (ER_WRONG_PARAMCOUNT_TO_NATIVE_FCT)

  Message: Incorrect parameter count in the call to native function '%s'

- Error: 1583 SQLSTATE: 42000 (ER_WRONG_PARAMETERS_TO_NATIVE_FCT)

  Message: Incorrect parameters in the call to native function '%s'

- Error: 1584 SQLSTATE: 42000 (ER_WRONG_PARAMETERS_TO_STORED_FCT)

  Message: Incorrect parameters in the call to stored function %s

- Error: 1585 SQLSTATE: HY000 (ER_NATIVE_FCT_NAME_COLLISION)

  Message: This function '%s' has the same name as a native function

- Error: 1586 SQLSTATE: 23000 (ER_DUP_ENTRY_WITH_KEY_NAME)

  Message: Duplicate entry '%s' for key '%s'

  The format string for this error is also used with ER_DUP_ENTRY.

- Error: 1587 SQLSTATE: HY000 (ER_BINLOG_PURGE_EMFILE)

Message: Too many files opened, please execute the command again

- Error: 1588 SQLSTATE: HY000 (ER_EVENT_CANNOT_CREATE_IN_THE_PAST)

  Message: Event execution time is in the past and ON COMPLETION NOT PRESERVE is set. The event was dropped immediately after creation.

- Error: 1589 SQLSTATE: HY000 (ER_EVENT_CANNOT_ALTER_IN_THE_PAST)

  Message: Event execution time is in the past and ON COMPLETION NOT PRESERVE is set. The event was not changed. Specify a time in the future.

- Error: 1590 SQLSTATE: HY000 (ER_SLAVE_INCIDENT)

  Message: The incident %s occured on the master. Message: %s

- Error: 1591 SQLSTATE: HY000 (ER_NO_PARTITION_FOR_GIVEN_VALUE_SILENT)

  Message: Table has no partition for some existing values

- Error: 1592 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_STATEMENT)

  Message: Unsafe statement written to the binary log using statement format since BINLOG_FORMAT = STATEMENT. %s

- Error: 1593 SQLSTATE: HY000 (ER_SLAVE_FATAL_ERROR)

  Message: Fatal error: %s

- Error: 1594 SQLSTATE: HY000 (ER_SLAVE_RELAY_LOG_READ_FAILURE)

  Message: Relay log read failure: %s

- Error: 1595 SQLSTATE: HY000 (ER_SLAVE_RELAY_LOG_WRITE_FAILURE)

  Message: Relay log write failure: %s

- Error: 1596 SQLSTATE: HY000 (ER_SLAVE_CREATE_EVENT_FAILURE)

  Message: Failed to create %s

- Error: 1597 SQLSTATE: HY000 (ER_SLAVE_MASTER_COM_FAILURE)

  Message: Master command %s failed: %s

- Error: 1598 SQLSTATE: HY000 (ER_BINLOG_LOGGING_IMPOSSIBLE)

  Message: Binary logging not possible. Message: %s

- Error: 1599 SQLSTATE: HY000 (ER_VIEW_NO_CREATION_CTX)

  Message: View `%s`.`%s` has no creation context

- Error: 1600 SQLSTATE: HY000 (ER_VIEW_INVALID_CREATION_CTX)

  Message: Creation context of view `%s`.`%s' is invalid

- Error: 1602 SQLSTATE: HY000 (ER_TRG_CORRUPTED_FILE)

  Message: Corrupted TRG file for table `%s`.`%s`

- Error: 1603 SQLSTATE: HY000 (ER_TRG_NO_CREATION_CTX)

  Message: Triggers for table `%s`.`%s` have no creation context

- Error: 1604 SQLSTATE: HY000 (ER_TRG_INVALID_CREATION_CTX)

  Message: Trigger creation context of table `%s`.`%s` is invalid

- Error: 1605 SQLSTATE: HY000 (ER_EVENT_INVALID_CREATION_CTX)

  Message: Creation context of event `%s`.`%s` is invalid

- Error: 1606 SQLSTATE: HY000 (ER_TRG_CANT_OPEN_TABLE)

  Message: Cannot open table for trigger `%s`.`%s`

- Error: 1609 SQLSTATE: HY000 (ER_NO_FORMAT_DESCRIPTION_EVENT_BEFORE_BINLOG_STATEMENT)

  Message: The BINLOG statement of type `%s` was not preceded by a format description BINLOG statement.

- Error: 1610 SQLSTATE: HY000 (ER_SLAVE_CORRUPT_EVENT)

  Message: Corrupted replication event was detected

- Error: 1612 SQLSTATE: HY000 (ER_LOG_PURGE_NO_FILE)

  Message: Being purged log %s was not found

- Error: 1613 SQLSTATE: XA106 (ER_XA_RBTIMEOUT)

  Message: XA_RBTIMEOUT: Transaction branch was rolled back: took too long

- Error: 1614 SQLSTATE: XA102 (ER_XA_RBDEADLOCK)

  Message: XA_RBDEADLOCK: Transaction branch was rolled back: deadlock was detected

- Error: 1615 SQLSTATE: HY000 (ER_NEED_REPREPARE)

  Message: Prepared statement needs to be re-prepared

- Error: 1617 SQLSTATE: HY000 (WARN_NO_MASTER_INFO)

  Message: The master info structure does not exist

- Error: 1618 SQLSTATE: HY000 (WARN_OPTION_IGNORED)

  Message: <%s> option ignored

- Error: 1619 SQLSTATE: HY000 (ER_PLUGIN_DELETE_BUILTIN)

  Message: Built-in plugins cannot be deleted

- Error: 1620 SQLSTATE: HY000 (WARN_PLUGIN_BUSY)

  Message: Plugin is busy and will be uninstalled on shutdown

- Error: 1621 SQLSTATE: HY000 (ER_VARIABLE_IS_READONLY)

  Message: %s variable '%s' is read-only. Use SET %s to assign the value

- Error: 1622 SQLSTATE: HY000 (ER_WARN_ENGINE_TRANSACTION_ROLLBACK)

  Message: Storage engine %s does not support rollback for this statement. Transaction rolled back and must be restarted

- Error: 1623 SQLSTATE: HY000 (ER_SLAVE_HEARTBEAT_FAILURE)

Message: Unexpected master's heartbeat data: %s

- Error: 1624 SQLSTATE: HY000 (ER_SLAVE_HEARTBEAT_VALUE_OUT_OF_RANGE)

  Message: The requested value for the heartbeat period is either negative or exceeds the maximum allowed (%s seconds).

- Error: 1625 SQLSTATE: HY000 (ER_NDB_REPLICATION_SCHEMA_ERROR)

  Message: Bad schema for mysql.ndb_replication table. Message: %s

- Error: 1626 SQLSTATE: HY000 (ER_CONFLICT_FN_PARSE_ERROR)

  Message: Error in parsing conflict function. Message: %s

- Error: 1627 SQLSTATE: HY000 (ER_EXCEPTIONS_WRITE_ERROR)

  Message: Write to exceptions table failed. Message: %s

- Error: 1628 SQLSTATE: HY000 (ER_TOO_LONG_TABLE_COMMENT)

  Message: Comment for table '%s' is too long (max = %lu)

- Error: 1629 SQLSTATE: HY000 (ER_TOO_LONG_FIELD_COMMENT)

  Message: Comment for field '%s' is too long (max = %lu)

- Error: 1630 SQLSTATE: 42000 (ER_FUNC_INEXISTENT_NAME_COLLISION)

  Message: FUNCTION %s does not exist. Check the 'Function Name Parsing and Resolution' section in the Reference Manual

- Error: 1631 SQLSTATE: HY000 (ER_DATABASE_NAME)

  Message: Database

- Error: 1632 SQLSTATE: HY000 (ER_TABLE_NAME)

  Message: Table

- Error: 1633 SQLSTATE: HY000 (ER_PARTITION_NAME)

  Message: Partition

- Error: 1634 SQLSTATE: HY000 (ER_SUBPARTITION_NAME)

  Message: Subpartition

- Error: 1635 SQLSTATE: HY000 (ER_TEMPORARY_NAME)

  Message: Temporary

- Error: 1636 SQLSTATE: HY000 (ER_RENAMED_NAME)

  Message: Renamed

- Error: 1637 SQLSTATE: HY000 (ER_TOO_MANY_CONCURRENT_TRXS)

  Message: Too many active concurrent transactions

- Error: 1638 SQLSTATE: HY000 (WARN_NON_ASCII_SEPARATOR_NOT_IMPLEMENTED)

  Message: Non-ASCII separator arguments are not fully supported

- Error: 1639 SQLSTATE: HY000 (ER_DEBUG_SYNC_TIMEOUT)

  Message: debug sync point wait timed out

- Error: 1640 SQLSTATE: HY000 (ER_DEBUG_SYNC_HIT_LIMIT)

  Message: debug sync point hit limit reached

- Error: 1641 SQLSTATE: 42000 (ER_DUP_SIGNAL_SET)

  Message: Duplicate condition information item '%s'

- Error: 1642 SQLSTATE: 01000 (ER_SIGNAL_WARN)

  Message: Unhandled user-defined warning condition

- Error: 1643 SQLSTATE: 02000 (ER_SIGNAL_NOT_FOUND)

  Message: Unhandled user-defined not found condition

- Error: 1644 SQLSTATE: HY000 (ER_SIGNAL_EXCEPTION)

  Message: Unhandled user-defined exception condition

- Error: 1645 SQLSTATE: 0K000 (ER_RESIGNAL_WITHOUT_ACTIVE_HANDLER)

  Message: RESIGNAL when handler not active

- Error: 1646 SQLSTATE: HY000 (ER_SIGNAL_BAD_CONDITION_TYPE)

  Message: SIGNAL/RESIGNAL can only use a CONDITION defined with SQLSTATE

- Error: 1647 SQLSTATE: HY000 (WARN_COND_ITEM_TRUNCATED)

  Message: Data truncated for condition item '%s'

- Error: 1648 SQLSTATE: HY000 (ER_COND_ITEM_TOO_LONG)

  Message: Data too long for condition item '%s'

- Error: 1649 SQLSTATE: HY000 (ER_UNKNOWN_LOCALE)

  Message: Unknown locale: '%s'

- Error: 1650 SQLSTATE: HY000 (ER_SLAVE_IGNORE_SERVER_IDS)

  Message: The requested server id %d clashes with the slave startup option --replicate-same-server-id

- Error: 1651 SQLSTATE: HY000 (ER_QUERY_CACHE_DISABLED)

  Message: Query cache is disabled; restart the server with query_cache_type=1 to enable it

- Error: 1652 SQLSTATE: HY000 (ER_SAME_NAME_PARTITION_FIELD)

  Message: Duplicate partition field name '%s'

- Error: 1653 SQLSTATE: HY000 (ER_PARTITION_COLUMN_LIST_ERROR)

  Message: Inconsistency in usage of column lists for partitioning

- Error: 1654 SQLSTATE: HY000 (ER_WRONG_TYPE_COLUMN_VALUE_ERROR)

  Message: Partition column values of incorrect type

- Error: 1655 SQLSTATE: HY000 (ER_TOO_MANY_PARTITION_FUNC_FIELDS_ERROR)

  Message: Too many fields in '%s'

- Error: 1656 SQLSTATE: HY000 (ER_MAXVALUE_IN_VALUES_IN)

  Message: Cannot use MAXVALUE as value in VALUES IN

- Error: 1657 SQLSTATE: HY000 (ER_TOO_MANY_VALUES_ERROR)

  Message: Cannot have more than one value for this type of %s partitioning

- Error: 1658 SQLSTATE: HY000 (ER_ROW_SINGLE_PARTITION_FIELD_ERROR)

  Message: Row expressions in VALUES IN only allowed for multi-field column partitioning

- Error: 1659 SQLSTATE: HY000 (ER_FIELD_TYPE_NOT_ALLOWED_AS_PARTITION_FIELD)

  Message: Field '%s' is of a not allowed type for this type of partitioning

- Error: 1660 SQLSTATE: HY000 (ER_PARTITION_FIELDS_TOO_LONG)

  Message: The total length of the partitioning fields is too large

- Error: 1661 SQLSTATE: HY000 (ER_BINLOG_ROW_ENGINE_AND_STMT_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since both row-incapable engines and statement-incapable engines are involved.

- Error: 1662 SQLSTATE: HY000 (ER_BINLOG_ROW_MODE_AND_STMT_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since BINLOG_FORMAT = ROW and at least one table uses a storage engine limited to statement-based logging.

- Error: 1663 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_AND_STMT_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since statement is unsafe, storage engine is limited to statement-based logging, and BINLOG_FORMAT = MIXED. %s

- Error: 1664 SQLSTATE: HY000 (ER_BINLOG_ROW_INJECTION_AND_STMT_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since statement is in row format and at least one table uses a storage engine limited to statement-based logging.

- Error: 1665 SQLSTATE: HY000 (ER_BINLOG_STMT_MODE_AND_ROW_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since BINLOG_FORMAT = STATEMENT and at least one table uses a storage engine limited to row-based logging.%s

- Error: 1666 SQLSTATE: HY000 (ER_BINLOG_ROW_INJECTION_AND_STMT_MODE)

  Message: Cannot execute statement: impossible to write to binary log since statement is in row format and BINLOG_FORMAT = STATEMENT.

- Error: 1667 SQLSTATE: HY000 (ER_BINLOG_MULTIPLE_ENGINES_AND_SELF_LOGGING_ENGINE)

  Message: Cannot execute statement: impossible to write to binary log since more than one engine is involved and at least one engine is self-logging.

- Error: 1668 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_LIMIT)

  Message: The statement is unsafe because it uses a LIMIT clause. This is unsafe because the set of rows included cannot be predicted.

- Error: 1670 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_SYSTEM_TABLE)

  Message: The statement is unsafe because it uses the general log, slow query log, or performance_schema table(s). This is unsafe because system tables may differ on slaves.

- Error: 1671 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_AUTOINC_COLUMNS)

  Message: Statement is unsafe because it invokes a trigger or a stored function that inserts into an AUTO_INCREMENT column. Inserted values cannot be logged correctly.

- Error: 1672 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_UDF)

  Message: Statement is unsafe because it uses a UDF which may not return the same value on the slave.

- Error: 1673 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_SYSTEM_VARIABLE)

  Message: Statement is unsafe because it uses a system variable that may have a different value on the slave.

- Error: 1674 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_SYSTEM_FUNCTION)

  Message: Statement is unsafe because it uses a system function that may return a different value on the slave.

- Error: 1675 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_NONTRANS_AFTER_TRANS)

  Message: Statement is unsafe because it accesses a non-transactional table after accessing a transactional table within the same transaction.

- Error: 1676 SQLSTATE: HY000 (ER_MESSAGE_AND_STATEMENT)

  Message: %s Statement: %s

- Error: 1677 SQLSTATE: HY000 (ER_SLAVE_CONVERSION_FAILED)

  Message: Column %d of table '%s.%s' cannot be converted from type '%s' to type '%s'

- Error: 1678 SQLSTATE: HY000 (ER_SLAVE_CANT_CREATE_CONVERSION)

  Message: Can't create conversion table for table '%s.%s'

- Error: 1679 SQLSTATE: HY000 (ER_INSIDE_TRANSACTION_PREVENTS_SWITCH_BINLOG_FORMAT)

  Message: Cannot modify @@session.binlog_format inside a transaction

- Error: 1680 SQLSTATE: HY000 (ER_PATH_LENGTH)

  Message: The path specified for %s is too long.

- Error: 1681 SQLSTATE: HY000 (ER_WARN_DEPRECATED_SYNTAX_NO_REPLACEMENT)

  Message: '%s' is deprecated and will be removed in a future release.

- Error: 1682 SQLSTATE: HY000 (ER_WRONG_NATIVE_TABLE_STRUCTURE)

  Message: Native table '%s'.'%s' has the wrong structure

- Error: 1683 SQLSTATE: HY000 (ER_WRONG_PERFSCHEMA_USAGE)

  Message: Invalid performance_schema usage.

- Error: 1684 SQLSTATE: HY000 (ER_WARN_I_S_SKIPPED_TABLE)

Message: Table '%s'.'%s' was skipped since its definition is being modified by concurrent DDL statement

- Error: 1685 SQLSTATE: HY000 (ER_INSIDE_TRANSACTION_PREVENTS_SWITCH_BINLOG_DIRECT)

  Message: Cannot modify @@session.binlog_direct_non_transactional_updates inside a transaction

- Error: 1686 SQLSTATE: HY000 (ER_STORED_FUNCTION_PREVENTS_SWITCH_BINLOG_DIRECT)

  Message: Cannot change the binlog direct flag inside a stored function or trigger

- Error: 1687 SQLSTATE: 42000 (ER_SPATIAL_MUST_HAVE_GEOM_COL)

  Message: A SPATIAL index may only contain a geometrical type column

- Error: 1688 SQLSTATE: HY000 (ER_TOO_LONG_INDEX_COMMENT)

  Message: Comment for index '%s' is too long (max = %lu)

- Error: 1689 SQLSTATE: HY000 (ER_LOCK_ABORTED)

  Message: Wait on a lock was aborted due to a pending exclusive lock

- Error: 1690 SQLSTATE: 22003 (ER_DATA_OUT_OF_RANGE)

  Message: %s value is out of range in '%s'

- Error: 1691 SQLSTATE: HY000 (ER_WRONG_SPVAR_TYPE_IN_LIMIT)

  Message: A variable of a non-integer based type in LIMIT clause

- Error: 1692 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_MULTIPLE_ENGINES_AND_SELF_LOGGING_ENGINE)

  Message: Mixing self-logging and non-self-logging engines in a statement is unsafe.

- Error: 1693 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_MIXED_STATEMENT)

  Message: Statement accesses nontransactional table as well as transactional or temporary table, and writes to any of them.

- Error: 1694 SQLSTATE: HY000 (ER_INSIDE_TRANSACTION_PREVENTS_SWITCH_SQL_LOG_BIN)

  Message: Cannot modify @@session.sql_log_bin inside a transaction

- Error: 1695 SQLSTATE: HY000 (ER_STORED_FUNCTION_PREVENTS_SWITCH_SQL_LOG_BIN)

  Message: Cannot change the sql_log_bin inside a stored function or trigger

- Error: 1696 SQLSTATE: HY000 (ER_FAILED_READ_FROM_PAR_FILE)

  Message: Failed to read from the .par file

- Error: 1697 SQLSTATE: HY000 (ER_VALUES_IS_NOT_INT_TYPE_ERROR)

  Message: VALUES value for partition '%s' must have type INT

- Error: 1698 SQLSTATE: 28000 (ER_ACCESS_DENIED_NO_PASSWORD_ERROR)

  Message: Access denied for user '%s'@'%s'

- Error: 1699 SQLSTATE: HY000 (ER_SET_PASSWORD_AUTH_PLUGIN)

Message: SET PASSWORD has no significance for users authenticating via plugins

- Error: 1701 SQLSTATE: 42000 (ER_TRUNCATE_ILLEGAL_FK)

  Message: Cannot truncate a table referenced in a foreign key constraint (%s)

- Error: 1702 SQLSTATE: HY000 (ER_PLUGIN_IS_PERMANENT)

  Message: Plugin '%s' is force_plus_permanent and can not be unloaded

- Error: 1703 SQLSTATE: HY000 (ER_SLAVE_HEARTBEAT_VALUE_OUT_OF_RANGE_MIN)

  Message: The requested value for the heartbeat period is less than 1 millisecond. The value is reset to 0, meaning that heartbeating will effectively be disabled.

- Error: 1704 SQLSTATE: HY000 (ER_SLAVE_HEARTBEAT_VALUE_OUT_OF_RANGE_MAX)

  Message: The requested value for the heartbeat period exceeds the value of `slave_net_timeout' seconds. A sensible value for the period should be less than the timeout.

- Error: 1705 SQLSTATE: HY000 (ER_STMT_CACHE_FULL)

  Message: Multi-row statements required more than 'max_binlog_stmt_cache_size' bytes of storage; increase this mysqld variable and try again

- Error: 1706 SQLSTATE: HY000 (ER_MULTI_UPDATE_KEY_CONFLICT)

  Message: Primary key/partition key update is not allowed since the table is updated both as '%s' and '%s'.

- Error: 1707 SQLSTATE: HY000 (ER_TABLE_NEEDS_REBUILD)

  Message: Table rebuild required. Please do "ALTER TABLE `%s` FORCE" or dump/reload to fix it!

- Error: 1708 SQLSTATE: HY000 (WARN_OPTION_BELOW_LIMIT)

  Message: The value of '%s' should be no less than the value of '%s'

- Error: 1709 SQLSTATE: HY000 (ER_INDEX_COLUMN_TOO_LONG)

  Message: Index column size too large. The maximum column size is %lu bytes.

- Error: 1710 SQLSTATE: HY000 (ER_ERROR_IN_TRIGGER_BODY)

  Message: Trigger '%s' has an error in its body: '%s'

- Error: 1711 SQLSTATE: HY000 (ER_ERROR_IN_UNKNOWN_TRIGGER_BODY)

  Message: Unknown trigger has an error in its body: '%s'

- Error: 1712 SQLSTATE: HY000 (ER_INDEX_CORRUPT)

  Message: Index %s is corrupted

- Error: 1713 SQLSTATE: HY000 (ER_UNDO_RECORD_TOO_BIG)

  Message: Undo log record is too big.

- Error: 1714 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_INSERT_IGNORE_SELECT)

  Message: INSERT IGNORE... SELECT is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are ignored. This order cannot be predicted and may differ on master and the slave.

- Error: 1715 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_INSERT_SELECT_UPDATE)

  Message: INSERT... SELECT... ON DUPLICATE KEY UPDATE is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are updated. This order cannot be predicted and may differ on master and the slave.

- Error: 1716 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_REPLACE_SELECT)

  Message: REPLACE... SELECT is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are replaced. This order cannot be predicted and may differ on master and the slave.

- Error: 1717 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_CREATE_IGNORE_SELECT)

  Message: CREATE... IGNORE SELECT is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are ignored. This order cannot be predicted and may differ on master and the slave.

- Error: 1718 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_CREATE_REPLACE_SELECT)

  Message: CREATE... REPLACE SELECT is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are replaced. This order cannot be predicted and may differ on master and the slave.

- Error: 1719 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_UPDATE_IGNORE)

  Message: UPDATE IGNORE is unsafe because the order in which rows are updated determines which (if any) rows are ignored. This order cannot be predicted and may differ on master and the slave.

- Error: 1720 SQLSTATE: HY000 (ER_PLUGIN_NO_UNINSTALL)

  Message: Plugin '%s' is marked as not dynamically uninstallable. You have to stop the server to uninstall it.

- Error: 1721 SQLSTATE: HY000 (ER_PLUGIN_NO_INSTALL)

  Message: Plugin '%s' is marked as not dynamically installable. You have to stop the server to install it.

- Error: 1722 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_WRITE_AUTOINC_SELECT)

  Message: Statements writing to a table with an auto-increment column after selecting from another table are unsafe because the order in which rows are retrieved determines what (if any) rows will be written. This order cannot be predicted and may differ on master and the slave.

- Error: 1723 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_CREATE_SELECT_AUTOINC)

  Message: CREATE TABLE... SELECT... on a table with an auto-increment column is unsafe because the order in which rows are retrieved by the SELECT determines which (if any) rows are inserted. This order cannot be predicted and may differ on master and the slave.

- Error: 1724 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_INSERT_TWO_KEYS)

  Message: INSERT... ON DUPLICATE KEY UPDATE on a table with more than one UNIQUE KEY is unsafe

- Error: 1725 SQLSTATE: HY000 (ER_TABLE_IN_FK_CHECK)

  Message: Table is being used in foreign key check.

- Error: 1726 SQLSTATE: HY000 (ER_UNSUPPORTED_ENGINE)

Message: Storage engine '%s' does not support system tables. [%s.%s]

- Error: 1727 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_AUTOINC_NOT_FIRST)

  Message: INSERT into autoincrement field which is not the first part in the composed primary key is unsafe.

- Error: 1728 SQLSTATE: HY000 (ER_CANNOT_LOAD_FROM_TABLE_V2)

  Message: Cannot load from %s.%s. The table is probably corrupted

- Error: 1729 SQLSTATE: HY000 (ER_MASTER_DELAY_VALUE_OUT_OF_RANGE)

  Message: The requested value %s for the master delay exceeds the maximum %u

- Error: 1730 SQLSTATE: HY000 (ER_ONLY_FD_AND_RBR_EVENTS_ALLOWED_IN_BINLOG_STATEMENT)

  Message: Only Format_description_log_event and row events are allowed in BINLOG statements (but %s was provided)

- Error: 1731 SQLSTATE: HY000 (ER_PARTITION_EXCHANGE_DIFFERENT_OPTION)

  Message: Non matching attribute '%s' between partition and table

- Error: 1732 SQLSTATE: HY000 (ER_PARTITION_EXCHANGE_PART_TABLE)

  Message: Table to exchange with partition is partitioned: '%s'

- Error: 1733 SQLSTATE: HY000 (ER_PARTITION_EXCHANGE_TEMP_TABLE)

  Message: Table to exchange with partition is temporary: '%s'

- Error: 1734 SQLSTATE: HY000 (ER_PARTITION_INSTEAD_OF_SUBPARTITION)

  Message: Subpartitioned table, use subpartition instead of partition

- Error: 1735 SQLSTATE: HY000 (ER_UNKNOWN_PARTITION)

  Message: Unknown partition '%s' in table '%s'

- Error: 1736 SQLSTATE: HY000 (ER_TABLES_DIFFERENT_METADATA)

  Message: Tables have different definitions

- Error: 1737 SQLSTATE: HY000 (ER_ROW_DOES_NOT_MATCH_PARTITION)

  Message: Found a row that does not match the partition

- Error: 1738 SQLSTATE: HY000 (ER_BINLOG_CACHE_SIZE_GREATER_THAN_MAX)

  Message: Option binlog_cache_size (%lu) is greater than max_binlog_cache_size (%lu); setting binlog_cache_size equal to max_binlog_cache_size.

- Error: 1739 SQLSTATE: HY000 (ER_WARN_INDEX_NOT_APPLICABLE)

  Message: Cannot use %s access on index '%s' due to type or collation conversion on field '%s'

- Error: 1740 SQLSTATE: HY000 (ER_PARTITION_EXCHANGE_FOREIGN_KEY)

  Message: Table to exchange with partition has foreign key references: '%s'

- Error: 1742 SQLSTATE: HY000 (ER_RPL_INFO_DATA_TOO_LONG)

Message: Data for column '%s' too long

- Error: 1743 SQLSTATE: HY000 (ER_NETWORK_READ_EVENT_CHECKSUM_FAILURE)

  Message: Replication event checksum verification failed while reading from network.

- Error: 1745 SQLSTATE: HY000 (ER_BINLOG_STMT_CACHE_SIZE_GREATER_THAN_MAX)

  Message: Option binlog_stmt_cache_size (%lu) is greater than max_binlog_stmt_cache_size (%lu); setting binlog_stmt_cache_size equal to max_binlog_stmt_cache_size.

- Error: 1746 SQLSTATE: HY000 (ER_CANT_UPDATE_TABLE_IN_CREATE_TABLE_SELECT)

  Message: Can't update table '%s' while '%s' is being created.

- Error: 1747 SQLSTATE: HY000 (ER_PARTITION_CLAUSE_ON_NONPARTITIONED)

  Message: PARTITION () clause on non partitioned table

- Error: 1748 SQLSTATE: HY000 (ER_ROW_DOES_NOT_MATCH_GIVEN_PARTITION_SET)

  Message: Found a row not matching the given partition set

- Error: 1750 SQLSTATE: HY000 (ER_CHANGE_RPL_INFO_REPOSITORY_FAILURE)

  Message: Failure while changing the type of replication repository: %s.

- Error: 1751 SQLSTATE: HY000 (ER_WARNING_NOT_COMPLETE_ROLLBACK_WITH_CREATED_TEMP_TABLE)

  Message: The creation of some temporary tables could not be rolled back.

- Error: 1752 SQLSTATE: HY000 (ER_WARNING_NOT_COMPLETE_ROLLBACK_WITH_DROPPED_TEMP_TABLE)

  Message: Some temporary tables were dropped, but these operations could not be rolled back.

- Error: 1753 SQLSTATE: HY000 (ER_MTS_FEATURE_IS_NOT_SUPPORTED)

  Message: %s is not supported in multi-threaded slave mode. %s

- Error: 1754 SQLSTATE: HY000 (ER_MTS_UPDATED_DBS_GREATER_MAX)

  Message: The number of modified databases exceeds the maximum %d; the database names will not be included in the replication event metadata.

- Error: 1755 SQLSTATE: HY000 (ER_MTS_CANT_PARALLEL)

  Message: Cannot execute the current event group in the parallel mode. Encountered event %s, relay-log name %s, position %s which prevents execution of this event group in parallel mode. Reason: %s.

- Error: 1756 SQLSTATE: HY000 (ER_MTS_INCONSISTENT_DATA)

  Message: %s

- Error: 1757 SQLSTATE: HY000 (ER_FULLTEXT_NOT_SUPPORTED_WITH_PARTITIONING)

  Message: FULLTEXT index is not supported for partitioned tables.

- Error: 1758 SQLSTATE: 35000 (ER_DA_INVALID_CONDITION_NUMBER)

  Message: Invalid condition number

- Error: 1759 SQLSTATE: HY000 (ER_INSECURE_PLAIN_TEXT)

  Message: Sending passwords in plain text without SSL/TLS is extremely insecure.

- Error: 1760 SQLSTATE: HY000 (ER_INSECURE_CHANGE_MASTER)

  Message: Storing MySQL user name or password information in the master info repository is not secure and is therefore not recommended. Please consider using the USER and PASSWORD connection options for START SLAVE; see the 'START SLAVE Syntax' in the MySQL Manual for more information.

- Error: 1761 SQLSTATE: 23000 (ER_FOREIGN_DUPLICATE_KEY_WITH_CHILD_INFO)

  Message: Foreign key constraint for table '%s', record '%s' would lead to a duplicate entry in table '%s', key '%s'

- Error: 1762 SQLSTATE: 23000 (ER_FOREIGN_DUPLICATE_KEY_WITHOUT_CHILD_INFO)

  Message: Foreign key constraint for table '%s', record '%s' would lead to a duplicate entry in a child table

- Error: 1763 SQLSTATE: HY000 (ER_SQLTHREAD_WITH_SECURE_SLAVE)

  Message: Setting authentication options is not possible when only the Slave SQL Thread is being started.

- Error: 1764 SQLSTATE: HY000 (ER_TABLE_HAS_NO_FT)

  Message: The table does not have FULLTEXT index to support this query

- Error: 1765 SQLSTATE: HY000 (ER_VARIABLE_NOT_SETTABLE_IN_SF_OR_TRIGGER)

  Message: The system variable %s cannot be set in stored functions or triggers.

- Error: 1766 SQLSTATE: HY000 (ER_VARIABLE_NOT_SETTABLE_IN_TRANSACTION)

  Message: The system variable %s cannot be set when there is an ongoing transaction.

- Error: 1769 SQLSTATE: HY000 (ER_SET_STATEMENT_CANNOT_INVOKE_FUNCTION)

  Message: The statement 'SET %s' cannot invoke a stored function.

- Error: 1770 SQLSTATE: HY000 (ER_GTID_NEXT_CANT_BE_AUTOMATIC_IF_GTID_NEXT_LIST_IS_NON_NULL)

  Message: The system variable @@SESSION.GTID_NEXT cannot be 'AUTOMATIC' when @@SESSION.GTID_NEXT_LIST is non-NULL.

- Error: 1772 SQLSTATE: HY000 (ER_MALFORMED_GTID_SET_SPECIFICATION)

  Message: Malformed GTID set specification '%s'.

- Error: 1773 SQLSTATE: HY000 (ER_MALFORMED_GTID_SET_ENCODING)

  Message: Malformed GTID set encoding.

- Error: 1774 SQLSTATE: HY000 (ER_MALFORMED_GTID_SPECIFICATION)

  Message: Malformed GTID specification '%s'.

- Error: 1775 SQLSTATE: HY000 (ER_GNO_EXHAUSTED)

  Message: Impossible to generate Global Transaction Identifier: the integer component reached the maximal value. Restart the server with a new server_uuid.

- Error: 1776 SQLSTATE: HY000 (ER_BAD_SLAVE_AUTO_POSITION)

  Message: Parameters MASTER_LOG_FILE, MASTER_LOG_POS, RELAY_LOG_FILE and RELAY_LOG_POS cannot be set when MASTER_AUTO_POSITION is active.

- Error: 1777 SQLSTATE: HY000 (ER_AUTO_POSITION_REQUIRES_GTID_MODE_NOT_OFF)

  Message: CHANGE MASTER TO MASTER_AUTO_POSITION = 1 cannot be executed because @@GLOBAL.GTID_MODE = OFF.

- Error: 1778 SQLSTATE: HY000 (ER_CANT_DO_IMPLICIT_COMMIT_IN_TRX_WHEN_GTID_NEXT_IS_SET)

  Message: Cannot execute statements with implicit commit inside a transaction when @@SESSION.GTID_NEXT == 'UUID:NUMBER'.

- Error: 1779 SQLSTATE: HY000 (ER_GTID_MODE_ON_REQUIRES_ENFORCE_GTID_CONSISTENCY_ON)

  Message: GTID_MODE = ON requires ENFORCE_GTID_CONSISTENCY = ON.

- Error: 1781 SQLSTATE: HY000 (ER_CANT_SET_GTID_NEXT_TO_GTID_WHEN_GTID_MODE_IS_OFF)

  Message: @@SESSION.GTID_NEXT cannot be set to UUID:NUMBER when @@GLOBAL.GTID_MODE = OFF.

- Error: 1782 SQLSTATE: HY000 (ER_CANT_SET_GTID_NEXT_TO_ANONYMOUS_WHEN_GTID_MODE_IS_ON)

  Message: @@SESSION.GTID_NEXT cannot be set to ANONYMOUS when @@GLOBAL.GTID_MODE = ON.

- Error: 1783 SQLSTATE: HY000 (ER_CANT_SET_GTID_NEXT_LIST_TO_NON_NULL_WHEN_GTID_MODE_IS_OFF)

  Message: @@SESSION.GTID_NEXT_LIST cannot be set to a non-NULL value when @@GLOBAL.GTID_MODE = OFF.

- Error: 1785 SQLSTATE: HY000 (ER_GTID_UNSAFE_NON_TRANSACTIONAL_TABLE)

  Message: Statement violates GTID consistency: Updates to non-transactional tables can only be done in either autocommitted statements or single-statement transactions, and never in the same statement as updates to transactional tables.

- Error: 1786 SQLSTATE: HY000 (ER_GTID_UNSAFE_CREATE_SELECT)

  Message: Statement violates GTID consistency: CREATE TABLE ... SELECT.

- Error: 1787 SQLSTATE: HY000 (ER_GTID_UNSAFE_CREATE_DROP_TEMPORARY_TABLE_IN_TRANSACTION)

  Message: Statement violates GTID consistency: CREATE TEMPORARY TABLE and DROP TEMPORARY TABLE can only be executed outside transactional context. These statements are also not allowed in a function or trigger because functions and triggers are also considered to be multi-statement transactions.

- Error: 1788 SQLSTATE: HY000 (ER_GTID_MODE_CAN_ONLY_CHANGE_ONE_STEP_AT_A_TIME)

  Message: The value of @@GLOBAL.GTID_MODE can only be changed one step at a time: OFF <-> OFF_PERMISSIVE <-> ON_PERMISSIVE <-> ON. Also note that this value must be stepped up or down simultaneously on all servers. See the Manual for instructions.

- Error: 1789 SQLSTATE: HY000 (ER_MASTER_HAS_PURGED_REQUIRED_GTIDS)

  Message: The slave is connecting using CHANGE MASTER TO MASTER_AUTO_POSITION = 1, but the master has purged binary logs containing GTIDs that the slave requires.

- Error: 1790 SQLSTATE: HY000 (ER_CANT_SET_GTID_NEXT_WHEN_OWNING_GTID)

  Message: @@SESSION.GTID_NEXT cannot be changed by a client that owns a GTID. The client owns %s. Ownership is released on COMMIT or ROLLBACK.

- Error: 1791 SQLSTATE: HY000 (ER_UNKNOWN_EXPLAIN_FORMAT)

  Message: Unknown EXPLAIN format name: '%s'

- Error: 1792 SQLSTATE: 25006 (ER_CANT_EXECUTE_IN_READ_ONLY_TRANSACTION)

  Message: Cannot execute statement in a READ ONLY transaction.

- Error: 1793 SQLSTATE: HY000 (ER_TOO_LONG_TABLE_PARTITION_COMMENT)

  Message: Comment for table partition '%s' is too long (max = %lu)

- Error: 1794 SQLSTATE: HY000 (ER_SLAVE_CONFIGURATION)

  Message: Slave is not configured or failed to initialize properly. You must at least set --server-id to enable either a master or a slave. Additional error messages can be found in the MySQL error log.

- Error: 1795 SQLSTATE: HY000 (ER_INNODB_FT_LIMIT)

  Message: InnoDB presently supports one FULLTEXT index creation at a time

- Error: 1796 SQLSTATE: HY000 (ER_INNODB_NO_FT_TEMP_TABLE)

  Message: Cannot create FULLTEXT index on temporary InnoDB table

- Error: 1797 SQLSTATE: HY000 (ER_INNODB_FT_WRONG_DOCID_COLUMN)

  Message: Column '%s' is of wrong type for an InnoDB FULLTEXT index

- Error: 1798 SQLSTATE: HY000 (ER_INNODB_FT_WRONG_DOCID_INDEX)

  Message: Index '%s' is of wrong type for an InnoDB FULLTEXT index

- Error: 1799 SQLSTATE: HY000 (ER_INNODB_ONLINE_LOG_TOO_BIG)

  Message: Creating index '%s' required more than 'innodb_online_alter_log_max_size' bytes of modification log. Please try again.

- Error: 1800 SQLSTATE: HY000 (ER_UNKNOWN_ALTER_ALGORITHM)

  Message: Unknown ALGORITHM '%s'

- Error: 1801 SQLSTATE: HY000 (ER_UNKNOWN_ALTER_LOCK)

  Message: Unknown LOCK type '%s'

- Error: 1802 SQLSTATE: HY000 (ER_MTS_CHANGE_MASTER_CANT_RUN_WITH_GAPS)

  Message: CHANGE MASTER cannot be executed when the slave was stopped with an error or killed in MTS mode. Consider using RESET SLAVE or START SLAVE UNTIL.

- Error: 1803 SQLSTATE: HY000 (ER_MTS_RECOVERY_FAILURE)

  Message: Cannot recover after SLAVE errored out in parallel execution mode. Additional error messages can be found in the MySQL error log.

- Error: 1804 SQLSTATE: HY000 (ER_MTS_RESET_WORKERS)

  Message: Cannot clean up worker info tables. Additional error messages can be found in the MySQL error log.

- Error: 1805 SQLSTATE: HY000 (ER_COL_COUNT_DOESNT_MATCH_CORRUPTED_V2)

  Message: Column count of %s.%s is wrong. Expected %d, found %d. The table is probably corrupted

- Error: 1806 SQLSTATE: HY000 (ER_SLAVE_SILENT_RETRY_TRANSACTION)

  Message: Slave must silently retry current transaction

- Error: 1807 SQLSTATE: HY000 (ER_DISCARD_FK_CHECKS_RUNNING)

  Message: There is a foreign key check running on table '%s'. Cannot discard the table.

- Error: 1808 SQLSTATE: HY000 (ER_TABLE_SCHEMA_MISMATCH)

  Message: Schema mismatch (%s)

- Error: 1809 SQLSTATE: HY000 (ER_TABLE_IN_SYSTEM_TABLESPACE)

  Message: Table '%s' in system tablespace

- Error: 1810 SQLSTATE: HY000 (ER_IO_READ_ERROR)

  Message: IO Read error: (%lu, %s) %s

- Error: 1811 SQLSTATE: HY000 (ER_IO_WRITE_ERROR)

  Message: IO Write error: (%lu, %s) %s

- Error: 1812 SQLSTATE: HY000 (ER_TABLESPACE_MISSING)

  Message: Tablespace is missing for table %s.

- Error: 1813 SQLSTATE: HY000 (ER_TABLESPACE_EXISTS)

  Message: Tablespace '%s' exists.

- Error: 1814 SQLSTATE: HY000 (ER_TABLESPACE_DISCARDED)

  Message: Tablespace has been discarded for table '%s'

- Error: 1815 SQLSTATE: HY000 (ER_INTERNAL_ERROR)

  Message: Internal error: %s

- Error: 1816 SQLSTATE: HY000 (ER_INNODB_IMPORT_ERROR)

  Message: ALTER TABLE %s IMPORT TABLESPACE failed with error %lu : '%s'

- Error: 1817 SQLSTATE: HY000 (ER_INNODB_INDEX_CORRUPT)

  Message: Index corrupt: %s

- Error: 1818 SQLSTATE: HY000 (ER_INVALID_YEAR_COLUMN_LENGTH)

  Message: Supports only YEAR or YEAR(4) column.

- Error: 1819 SQLSTATE: HY000 (ER_NOT_VALID_PASSWORD)

  Message: Your password does not satisfy the current policy requirements

- Error: 1820 SQLSTATE: HY000 (ER_MUST_CHANGE_PASSWORD)

  Message: You must reset your password using ALTER USER statement before executing this statement.

- Error: 1821 SQLSTATE: HY000 (ER_FK_NO_INDEX_CHILD)

  Message: Failed to add the foreign key constaint. Missing index for constraint '%s' in the foreign table '%s'

- Error: 1822 SQLSTATE: HY000 (ER_FK_NO_INDEX_PARENT)

  Message: Failed to add the foreign key constaint. Missing index for constraint '%s' in the referenced table '%s'

- Error: 1823 SQLSTATE: HY000 (ER_FK_FAIL_ADD_SYSTEM)

  Message: Failed to add the foreign key constraint '%s' to system tables

- Error: 1824 SQLSTATE: HY000 (ER_FK_CANNOT_OPEN_PARENT)

  Message: Failed to open the referenced table '%s'

- Error: 1825 SQLSTATE: HY000 (ER_FK_INCORRECT_OPTION)

  Message: Failed to add the foreign key constraint on table '%s'. Incorrect options in FOREIGN KEY constraint '%s'

- Error: 1826 SQLSTATE: HY000 (ER_FK_DUP_NAME)

  Message: Duplicate foreign key constraint name '%s'

- Error: 1827 SQLSTATE: HY000 (ER_PASSWORD_FORMAT)

  Message: The password hash doesn't have the expected format. Check if the correct password algorithm is being used with the PASSWORD() function.

- Error: 1828 SQLSTATE: HY000 (ER_FK_COLUMN_CANNOT_DROP)

  Message: Cannot drop column '%s': needed in a foreign key constraint '%s'

- Error: 1829 SQLSTATE: HY000 (ER_FK_COLUMN_CANNOT_DROP_CHILD)

  Message: Cannot drop column '%s': needed in a foreign key constraint '%s' of table '%s'

- Error: 1830 SQLSTATE: HY000 (ER_FK_COLUMN_NOT_NULL)

  Message: Column '%s' cannot be NOT NULL: needed in a foreign key constraint '%s' SET NULL

- Error: 1831 SQLSTATE: HY000 (ER_DUP_INDEX)

  Message: Duplicate index '%s' defined on the table '%s.%s'. This is deprecated and will be disallowed in a future release.

- Error: 1832 SQLSTATE: HY000 (ER_FK_COLUMN_CANNOT_CHANGE)

  Message: Cannot change column '%s': used in a foreign key constraint '%s'

- Error: 1833 SQLSTATE: HY000 (ER_FK_COLUMN_CANNOT_CHANGE_CHILD)

  Message: Cannot change column '%s': used in a foreign key constraint '%s' of table '%s'

- Error: 1835 SQLSTATE: HY000 (ER_MALFORMED_PACKET)

Message: Malformed communication packet.

- Error: 1836 SQLSTATE: HY000 (ER_READ_ONLY_MODE)

  Message: Running in read-only mode

- Error: 1837 SQLSTATE: HY000 (ER_GTID_NEXT_TYPE_UNDEFINED_GROUP)

  Message: When @@SESSION.GTID_NEXT is set to a GTID, you must explicitly set it to a different value after a COMMIT or ROLLBACK. Please check GTID_NEXT variable manual page for detailed explanation. Current @@SESSION.GTID_NEXT is '%s'.

- Error: 1838 SQLSTATE: HY000 (ER_VARIABLE_NOT_SETTABLE_IN_SP)

  Message: The system variable %s cannot be set in stored procedures.

- Error: 1840 SQLSTATE: HY000 (ER_CANT_SET_GTID_PURGED_WHEN_GTID_EXECUTED_IS_NOT_EMPTY)

  Message: @@GLOBAL.GTID_PURGED can only be set when @@GLOBAL.GTID_EXECUTED is empty.

- Error: 1841 SQLSTATE: HY000 (ER_CANT_SET_GTID_PURGED_WHEN_OWNED_GTIDS_IS_NOT_EMPTY)

  Message: @@GLOBAL.GTID_PURGED can only be set when there are no ongoing transactions (not even in other clients).

- Error: 1842 SQLSTATE: HY000 (ER_GTID_PURGED_WAS_CHANGED)

  Message: @@GLOBAL.GTID_PURGED was changed from '%s' to '%s'.

- Error: 1843 SQLSTATE: HY000 (ER_GTID_EXECUTED_WAS_CHANGED)

  Message: @@GLOBAL.GTID_EXECUTED was changed from '%s' to '%s'.

- Error: 1844 SQLSTATE: HY000 (ER_BINLOG_STMT_MODE_AND_NO_REPL_TABLES)

  Message: Cannot execute statement: impossible to write to binary log since BINLOG_FORMAT = STATEMENT, and both replicated and non replicated tables are written to.

- Error: 1845 SQLSTATE: 0A000 (ER_ALTER_OPERATION_NOT_SUPPORTED)

  Message: %s is not supported for this operation. Try %s.

- Error: 1846 SQLSTATE: 0A000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON)

  Message: %s is not supported. Reason: %s. Try %s.

- Error: 1847 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_COPY)

  Message: COPY algorithm requires a lock

- Error: 1848 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_PARTITION)

  Message: Partition specific operations do not yet support LOCK/ALGORITHM

- Error: 1849 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_FK_RENAME)

  Message: Columns participating in a foreign key are renamed

- Error: 1850 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_COLUMN_TYPE)

  Message: Cannot change column type INPLACE

- Error: 1851 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_FK_CHECK)

  Message: Adding foreign keys needs foreign_key_checks=OFF

- Error: 1853 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_NOPK)

  Message: Dropping a primary key is not allowed without also adding a new primary key

- Error: 1854 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_AUTOINC)

  Message: Adding an auto-increment column requires a lock

- Error: 1855 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_HIDDEN_FTS)

  Message: Cannot replace hidden FTS_DOC_ID with a user-visible one

- Error: 1856 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_CHANGE_FTS)

  Message: Cannot drop or rename FTS_DOC_ID

- Error: 1857 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_FTS)

  Message: Fulltext index creation requires a lock

- Error: 1858 SQLSTATE: HY000 (ER_SQL_SLAVE_SKIP_COUNTER_NOT_SETTABLE_IN_GTID_MODE)

  Message: sql_slave_skip_counter can not be set when the server is running with @@GLOBAL.GTID_MODE = ON. Instead, for each transaction that you want to skip, generate an empty transaction with the same GTID as the transaction

- Error: 1859 SQLSTATE: 23000 (ER_DUP_UNKNOWN_IN_INDEX)

  Message: Duplicate entry for key '%s'

- Error: 1860 SQLSTATE: HY000 (ER_IDENT_CAUSES_TOO_LONG_PATH)

  Message: Long database name and identifier for object resulted in path length exceeding %d characters. Path: '%s'.

- Error: 1861 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_NOT_NULL)

  Message: cannot silently convert NULL values, as required in this SQL_MODE

- Error: 1862 SQLSTATE: HY000 (ER_MUST_CHANGE_PASSWORD_LOGIN)

  Message: Your password has expired. To log in you must change it using a client that supports expired passwords.

- Error: 1863 SQLSTATE: HY000 (ER_ROW_IN_WRONG_PARTITION)

  Message: Found a row in wrong partition %s

- Error: 1864 SQLSTATE: HY000 (ER_MTS_EVENT_BIGGER_PENDING_JOBS_SIZE_MAX)

  Message: Cannot schedule event %s, relay-log name %s, position %s to Worker thread because its size %lu exceeds %lu of slave_pending_jobs_size_max.

- Error: 1866 SQLSTATE: HY000 (ER_BINLOG_LOGICAL_CORRUPTION)

  Message: The binary log file '%s' is logically corrupted: %s

- Error: 1867 SQLSTATE: HY000 (ER_WARN_PURGE_LOG_IN_USE)

  Message: file %s was not purged because it was being read by %d thread(s), purged only %d out of %d files.

- Error: 1868 SQLSTATE: HY000 (ER_WARN_PURGE_LOG_IS_ACTIVE)

  Message: file %s was not purged because it is the active log file.

- Error: 1869 SQLSTATE: HY000 (ER_AUTO_INCREMENT_CONFLICT)

  Message: Auto-increment value in UPDATE conflicts with internally generated values

- Error: 1870 SQLSTATE: HY000 (WARN_ON_BLOCKHOLE_IN_RBR)

  Message: Row events are not logged for %s statements that modify BLACKHOLE tables in row format. Table(s): '%s'

- Error: 1871 SQLSTATE: HY000 (ER_SLAVE_MI_INIT_REPOSITORY)

  Message: Slave failed to initialize master info structure from the repository

- Error: 1872 SQLSTATE: HY000 (ER_SLAVE_RLI_INIT_REPOSITORY)

  Message: Slave failed to initialize relay log info structure from the repository

- Error: 1873 SQLSTATE: 28000 (ER_ACCESS_DENIED_CHANGE_USER_ERROR)

  Message: Access denied trying to change to user '%s'@'%s' (using password: %s). Disconnecting.

- Error: 1874 SQLSTATE: HY000 (ER_INNODB_READ_ONLY)

  Message: InnoDB is in read only mode.

- Error: 1875 SQLSTATE: HY000 (ER_STOP_SLAVE_SQL_THREAD_TIMEOUT)

  Message: STOP SLAVE command execution is incomplete: Slave SQL thread got the stop signal, thread is busy, SQL thread will stop once the current task is complete.

- Error: 1876 SQLSTATE: HY000 (ER_STOP_SLAVE_IO_THREAD_TIMEOUT)

  Message: STOP SLAVE command execution is incomplete: Slave IO thread got the stop signal, thread is busy, IO thread will stop once the current task is complete.

- Error: 1877 SQLSTATE: HY000 (ER_TABLE_CORRUPT)

  Message: Operation cannot be performed. The table '%s.%s' is missing, corrupt or contains bad data.

- Error: 1878 SQLSTATE: HY000 (ER_TEMP_FILE_WRITE_FAILURE)

  Message: Temporary file write failure.

- Error: 1879 SQLSTATE: HY000 (ER_INNODB_FT_AUX_NOT_HEX_ID)

  Message: Upgrade index name failed, please use create index(alter table) algorithm copy to rebuild index.

- Error: 1880 SQLSTATE: HY000 (ER_OLD_TEMPORALS_UPGRADED)

Message: TIME/TIMESTAMP/DATETIME columns of old format have been upgraded to the new format.

- Error: 1881 SQLSTATE: HY000 (ER_INNODB_FORCED_RECOVERY)

  Message: Operation not allowed when innodb_forced_recovery > 0.

- Error: 1882 SQLSTATE: HY000 (ER_AES_INVALID_IV)

  Message: The initialization vector supplied to %s is too short. Must be at least %d bytes long

- Error: 1883 SQLSTATE: HY000 (ER_PLUGIN_CANNOT_BE_UNINSTALLED)

  Message: Plugin '%s' cannot be uninstalled now. %s

- Error: 1884 SQLSTATE: HY000 (ER_GTID_UNSAFE_BINLOG_SPLITTABLE_STATEMENT_AND_GTID_GROUP)

  Message: Cannot execute statement because it needs to be written to the binary log as multiple statements, and this is not allowed when @@SESSION.GTID_NEXT == 'UUID:NUMBER'.

- Error: 1885 SQLSTATE: HY000 (ER_SLAVE_HAS_MORE_GTIDS_THAN_MASTER)

  Message: Slave has more GTIDs than the master has, using the master's SERVER_UUID. This may indicate that the end of the binary log was truncated or that the last binary log file was lost, e.g., after a power or disk failure when sync_binlog != 1. The master may or may not have rolled back transactions that were already replicated to the slave. Suggest to replicate any transactions that master has rolled back from slave to master, and/or commit empty transactions on master to account for transactions that have been committed on master but are not included in GTID_EXECUTED.

- Error: 3000 SQLSTATE: HY000 (ER_FILE_CORRUPT)

  Message: File %s is corrupted

- Error: 3001 SQLSTATE: HY000 (ER_ERROR_ON_MASTER)

  Message: Query partially completed on the master (error on master: %d) and was aborted. There is a chance that your master is inconsistent at this point. If you are sure that your master is ok, run this query manually on the slave and then restart the slave with SET GLOBAL SQL_SLAVE_SKIP_COUNTER=1; START SLAVE;. Query:'%s'

- Error: 3002 SQLSTATE: HY000 (ER_INCONSISTENT_ERROR)

  Message: Query caused different errors on master and slave. Error on master: message (format)='%s' error code=%d; Error on slave:actual message='%s', error code=%d. Default database:'%s'. Query:'%s'

- Error: 3003 SQLSTATE: HY000 (ER_STORAGE_ENGINE_NOT_LOADED)

  Message: Storage engine for table '%s'.'%s' is not loaded.

- Error: 3004 SQLSTATE: 0Z002 (ER_GET_STACKED_DA_WITHOUT_ACTIVE_HANDLER)

  Message: GET STACKED DIAGNOSTICS when handler not active

- Error: 3005 SQLSTATE: HY000 (ER_WARN_LEGACY_SYNTAX_CONVERTED)

  Message: %s is no longer supported. The statement was converted to %s.

- Error: 3006 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_FULLTEXT_PLUGIN)

  Message: Statement is unsafe because it uses a fulltext parser plugin which may not return the same value on the slave.

- Error: 3007 SQLSTATE: HY000 (ER_CANNOT_DISCARD_TEMPORARY_TABLE)

  Message: Cannot DISCARD/IMPORT tablespace associated with temporary table

- Error: 3008 SQLSTATE: HY000 (ER_FK_DEPTH_EXCEEDED)

  Message: Foreign key cascade delete/update exceeds max depth of %d.

- Error: 3009 SQLSTATE: HY000 (ER_COL_COUNT_DOESNT_MATCH_PLEASE_UPDATE_V2)

  Message: Column count of %s.%s is wrong. Expected %d, found %d. Created with MySQL %d, now running %d. Please use mysql_upgrade to fix this error.

- Error: 3010 SQLSTATE: HY000 (ER_WARN_TRIGGER_DOESNT_HAVE_CREATED)

  Message: Trigger %s.%s.%s does not have CREATED attribute.

- Error: 3011 SQLSTATE: HY000 (ER_REFERENCED_TRG_DOES_NOT_EXIST)

  Message: Referenced trigger '%s' for the given action time and event type does not exist.

- Error: 3012 SQLSTATE: HY000 (ER_EXPLAIN_NOT_SUPPORTED)

  Message: EXPLAIN FOR CONNECTION command is supported only for SELECT/UPDATE/INSERT/DELETE/REPLACE

- Error: 3013 SQLSTATE: HY000 (ER_INVALID_FIELD_SIZE)

  Message: Invalid size for column '%s'.

- Error: 3014 SQLSTATE: HY000 (ER_MISSING_HA_CREATE_OPTION)

  Message: Table storage engine '%s' found required create option missing

- Error: 3015 SQLSTATE: HY000 (ER_ENGINE_OUT_OF_MEMORY)

  Message: Out of memory in storage engine '%s'.

- Error: 3016 SQLSTATE: HY000 (ER_PASSWORD_EXPIRE_ANONYMOUS_USER)

  Message: The password for anonymous user cannot be expired.

- Error: 3017 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_MUST_STOP)

  Message: This operation cannot be performed with a running slave sql thread; run STOP SLAVE SQL_THREAD first

- Error: 3018 SQLSTATE: HY000 (ER_NO_FT_MATERIALIZED_SUBQUERY)

  Message: Cannot create FULLTEXT index on materialized subquery

- Error: 3019 SQLSTATE: HY000 (ER_INNODB_UNDO_LOG_FULL)

  Message: Undo Log error: %s

- Error: 3020 SQLSTATE: 2201E (ER_INVALID_ARGUMENT_FOR_LOGARITHM)

  Message: Invalid argument for logarithm

- Error: 3021 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_IO_THREAD_MUST_STOP)

  Message: This operation cannot be performed with a running slave io thread; run STOP SLAVE IO_THREAD FOR CHANNEL '%s' first.

- Error: 3022 SQLSTATE: HY000 (ER_WARN_OPEN_TEMP_TABLES_MUST_BE_ZERO)

  Message: This operation may not be safe when the slave has temporary tables. The tables will be kept open until the server restarts or until the tables are deleted by any replicated DROP statement. Suggest to wait until slave_open_temp_tables = 0.

- Error: 3023 SQLSTATE: HY000 (ER_WARN_ONLY_MASTER_LOG_FILE_NO_POS)

  Message: CHANGE MASTER TO with a MASTER_LOG_FILE clause but no MASTER_LOG_POS clause may not be safe. The old position value may not be valid for the new binary log file.

- Error: 3024 SQLSTATE: HY000 (ER_QUERY_TIMEOUT)

  Message: Query execution was interrupted, maximum statement execution time exceeded

- Error: 3025 SQLSTATE: HY000 (ER_NON_RO_SELECT_DISABLE_TIMER)

  Message: Select is not a read only statement, disabling timer

- Error: 3026 SQLSTATE: HY000 (ER_DUP_LIST_ENTRY)

  Message: Duplicate entry '%s'.

- Error: 3028 SQLSTATE: HY000 (ER_AGGREGATE_ORDER_FOR_UNION)

  Message: Expression #%u of ORDER BY contains aggregate function and applies to a UNION

- Error: 3029 SQLSTATE: HY000 (ER_AGGREGATE_ORDER_NON_AGG_QUERY)

  Message: Expression #%u of ORDER BY contains aggregate function and applies to the result of a non-aggregated query

- Error: 3030 SQLSTATE: HY000 (ER_SLAVE_WORKER_STOPPED_PREVIOUS_THD_ERROR)

  Message: Slave worker has stopped after at least one previous worker encountered an error when slave-preserve-commit-order was enabled. To preserve commit order, the last transaction executed by this thread has not been committed. When restarting the slave after fixing any failed threads, you should fix this worker as well.

- Error: 3031 SQLSTATE: HY000 (ER_DONT_SUPPORT_SLAVE_PRESERVE_COMMIT_ORDER)

  Message: slave_preserve_commit_order is not supported %s.

- Error: 3032 SQLSTATE: HY000 (ER_SERVER_OFFLINE_MODE)

  Message: The server is currently in offline mode

- Error: 3033 SQLSTATE: HY000 (ER_GIS_DIFFERENT_SRIDS)

  Message: Binary geometry function %s given two geometries of different srids: %u and %u, which should have been identical.

  Geometry values passed as arguments to spatial functions must have the same SRID value.

- Error: 3034 SQLSTATE: HY000 (ER_GIS_UNSUPPORTED_ARGUMENT)

  Message: Calling geometry function %s with unsupported types of arguments.

  A spatial function was called with a combination of argument types that the function does not support.

- Error: 3035 SQLSTATE: HY000 (ER_GIS_UNKNOWN_ERROR)

  Message: Unknown GIS error occured in function %s.

- Error: 3036 SQLSTATE: HY000 (ER_GIS_UNKNOWN_EXCEPTION)

  Message: Unknown exception caught in GIS function %s.

- Error: 3037 SQLSTATE: 22023 (ER_GIS_INVALID_DATA)

  Message: Invalid GIS data provided to function %s.

  A spatial function was called with an argument not recognized as a valid geometry value.

- Error: 3038 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_EMPTY_INPUT_EXCEPTION)

  Message: The geometry has no data in function %s.

- Error: 3039 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_CENTROID_EXCEPTION)

  Message: Unable to calculate centroid because geometry is empty in function %s.

- Error: 3040 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_OVERLAY_INVALID_INPUT_EXCEPTION)

  Message: Geometry overlay calculation error: geometry data is invalid in function %s.

- Error: 3041 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_TURN_INFO_EXCEPTION)

  Message: Geometry turn info calculation error: geometry data is invalid in function %s.

- Error: 3042 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_SELF_INTERSECTION_POINT_EXCEPTION)

  Message: Analysis procedures of intersection points interrupted unexpectedly in function %s.

- Error: 3043 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_UNKNOWN_EXCEPTION)

  Message: Unknown exception thrown in function %s.

- Error: 3044 SQLSTATE: HY000 (ER_STD_BAD_ALLOC_ERROR)

  Message: Memory allocation error: %s in function %s.

- Error: 3045 SQLSTATE: HY000 (ER_STD_DOMAIN_ERROR)

  Message: Domain error: %s in function %s.

- Error: 3046 SQLSTATE: HY000 (ER_STD_LENGTH_ERROR)

  Message: Length error: %s in function %s.

- Error: 3047 SQLSTATE: HY000 (ER_STD_INVALID_ARGUMENT)

  Message: Invalid argument error: %s in function %s.

- Error: 3048 SQLSTATE: HY000 (ER_STD_OUT_OF_RANGE_ERROR)

  Message: Out of range error: %s in function %s.

- Error: 3049 SQLSTATE: HY000 (ER_STD_OVERFLOW_ERROR)

  Message: Overflow error error: %s in function %s.

- Error: 3050 SQLSTATE: HY000 (ER_STD_RANGE_ERROR)

  Message: Range error: %s in function %s.

- Error: 3051 SQLSTATE: HY000 (ER_STD_UNDERFLOW_ERROR)

Message: Underflow error: %s in function %s.

- Error: 3052 SQLSTATE: HY000 (ER_STD_LOGIC_ERROR)

  Message: Logic error: %s in function %s.

- Error: 3053 SQLSTATE: HY000 (ER_STD_RUNTIME_ERROR)

  Message: Runtime error: %s in function %s.

- Error: 3054 SQLSTATE: HY000 (ER_STD_UNKNOWN_EXCEPTION)

  Message: Unknown exception: %s in function %s.

- Error: 3055 SQLSTATE: HY000 (ER_GIS_DATA_WRONG_ENDIANESS)

  Message: Geometry byte string must be little endian.

- Error: 3056 SQLSTATE: HY000 (ER_CHANGE_MASTER_PASSWORD_LENGTH)

  Message: The password provided for the replication user exceeds the maximum length of 32 characters

- Error: 3057 SQLSTATE: 42000 (ER_USER_LOCK_WRONG_NAME)

  Message: Incorrect user-level lock name '%s'.

- Error: 3058 SQLSTATE: HY000 (ER_USER_LOCK_DEADLOCK)

  Message: Deadlock found when trying to get user-level lock; try rolling back transaction/releasing locks and restarting lock acquisition.

  This error is returned when the metdata locking subsystem detects a deadlock for an attempt to acquire a named lock with GET_LOCK.

- Error: 3059 SQLSTATE: HY000 (ER_REPLACE_INACCESSIBLE_ROWS)

  Message: REPLACE cannot be executed as it requires deleting rows that are not in the view

- Error: 3060 SQLSTATE: HY000 (ER_ALTER_OPERATION_NOT_SUPPORTED_REASON_GIS)

  Message: Do not support online operation on table with GIS index

- Error: 3061 SQLSTATE: 42000 (ER_ILLEGAL_USER_VAR)

  Message: User variable name '%s' is illegal

- Error: 3062 SQLSTATE: HY000 (ER_GTID_MODE_OFF)

  Message: Cannot %s when GTID_MODE = OFF.

- Error: 3064 SQLSTATE: HY000 (ER_INCORRECT_TYPE)

  Message: Incorrect type for argument %s in function %s.

- Error: 3065 SQLSTATE: HY000 (ER_FIELD_IN_ORDER_NOT_SELECT)

  Message: Expression #%u of ORDER BY clause is not in SELECT list, references column '%s' which is not in SELECT list; this is incompatible with %s

- Error: 3066 SQLSTATE: HY000 (ER_AGGREGATE_IN_ORDER_NOT_SELECT)

  Message: Expression #%u of ORDER BY clause is not in SELECT list, contains aggregate function; this is incompatible with %s

- Error: 3067 SQLSTATE: HY000 (ER_INVALID_RPL_WILD_TABLE_FILTER_PATTERN)

  Message: Supplied filter list contains a value which is not in the required format 'db_pattern.table_pattern'

- Error: 3068 SQLSTATE: 08S01 (ER_NET_OK_PACKET_TOO_LARGE)

  Message: OK packet too large

- Error: 3069 SQLSTATE: HY000 (ER_INVALID_JSON_DATA)

  Message: Invalid JSON data provided to function %s: %s

- Error: 3070 SQLSTATE: HY000 (ER_INVALID_GEOJSON_MISSING_MEMBER)

  Message: Invalid GeoJSON data provided to function %s: Missing required member '%s'

- Error: 3071 SQLSTATE: HY000 (ER_INVALID_GEOJSON_WRONG_TYPE)

  Message: Invalid GeoJSON data provided to function %s: Member '%s' must be of type '%s'

- Error: 3072 SQLSTATE: HY000 (ER_INVALID_GEOJSON_UNSPECIFIED)

  Message: Invalid GeoJSON data provided to function %s

- Error: 3073 SQLSTATE: HY000 (ER_DIMENSION_UNSUPPORTED)

  Message: Unsupported number of coordinate dimensions in function %s: Found %u, expected %u

- Error: 3074 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_DOES_NOT_EXIST)

  Message: Slave channel '%s' does not exist.

- Error: 3076 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_NAME_INVALID_OR_TOO_LONG)

  Message: Couldn't create channel: Channel name is either invalid or too long.

- Error: 3077 SQLSTATE: HY000 (ER_SLAVE_NEW_CHANNEL_WRONG_REPOSITORY)

  Message: To have multiple channels, repository cannot be of type FILE; Please check the repository configuration and convert them to TABLE.

- Error: 3079 SQLSTATE: HY000 (ER_SLAVE_MULTIPLE_CHANNELS_CMD)

  Message: Multiple channels exist on the slave. Please provide channel name as an argument.

- Error: 3080 SQLSTATE: HY000 (ER_SLAVE_MAX_CHANNELS_EXCEEDED)

  Message: Maximum number of replication channels allowed exceeded.

- Error: 3081 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_MUST_STOP)

  Message: This operation cannot be performed with running replication threads; run STOP SLAVE FOR CHANNEL '%s' first

- Error: 3082 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_NOT_RUNNING)

  Message: This operation requires running replication threads; configure slave and run START SLAVE FOR CHANNEL '%s'

- Error: 3083 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_WAS_RUNNING)

  Message: Replication thread(s) for channel '%s' are already runnning.

- Error: 3084 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_WAS_NOT_RUNNING)

Message: Replication thread(s) for channel '%s' are already stopped.

- Error: 3085 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_SQL_THREAD_MUST_STOP)

  Message: This operation cannot be performed with a running slave sql thread; run STOP SLAVE SQL_THREAD FOR CHANNEL '%s' first.

- Error: 3086 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_SQL_SKIP_COUNTER)

  Message: When sql_slave_skip_counter > 0, it is not allowed to start more than one SQL thread by using 'START SLAVE [SQL_THREAD]'. Value of sql_slave_skip_counter can only be used by one SQL thread at a time. Please use 'START SLAVE [SQL_THREAD] FOR CHANNEL' to start the SQL thread which will use the value of sql_slave_skip_counter.

- Error: 3087 SQLSTATE: HY000 (ER_WRONG_FIELD_WITH_GROUP_V2)

  Message: Expression #%u of %s is not in GROUP BY clause and contains nonaggregated column '%s' which is not functionally dependent on columns in GROUP BY clause; this is incompatible with sql_mode=only_full_group_by

- Error: 3088 SQLSTATE: HY000 (ER_MIX_OF_GROUP_FUNC_AND_FIELDS_V2)

  Message: In aggregated query without GROUP BY, expression #%u of %s contains nonaggregated column '%s'; this is incompatible with sql_mode=only_full_group_by

- Error: 3089 SQLSTATE: HY000 (ER_WARN_DEPRECATED_SYSVAR_UPDATE)

  Message: Updating '%s' is deprecated. It will be made read-only in a future release.

- Error: 3090 SQLSTATE: HY000 (ER_WARN_DEPRECATED_SQLMODE)

  Message: Changing sql mode '%s' is deprecated. It will be removed in a future release.

- Error: 3091 SQLSTATE: HY000 (ER_CANNOT_LOG_PARTIAL_DROP_DATABASE_WITH_GTID)

  Message: DROP DATABASE failed; some tables may have been dropped but the database directory remains. The GTID has not been added to GTID_EXECUTED and the statement was not written to the binary log. Fix this as follows: (1) remove all files from the database directory %s; (2) SET GTID_NEXT='%s'; (3) DROP DATABASE `%s`.

- Error: 3092 SQLSTATE: HY000 (ER_GROUP_REPLICATION_CONFIGURATION)

  Message: The server is not configured properly to be an active member of the group. Please see more details on error log.

- Error: 3093 SQLSTATE: HY000 (ER_GROUP_REPLICATION_RUNNING)

  Message: The START GROUP_REPLICATION command failed since the group is already running.

- Error: 3094 SQLSTATE: HY000 (ER_GROUP_REPLICATION_APPLIER_INIT_ERROR)

  Message: The START GROUP_REPLICATION command failed as the applier module failed to start.

- Error: 3095 SQLSTATE: HY000 (ER_GROUP_REPLICATION_STOP_APPLIER_THREAD_TIMEOUT)

  Message: The STOP GROUP_REPLICATION command execution is incomplete: The applier thread got the stop signal while it was busy. The applier thread will stop once the current task is complete.

- Error: 3096 SQLSTATE: HY000 (ER_GROUP_REPLICATION_COMMUNICATION_LAYER_SESSION_ERROR)

  Message: The START GROUP_REPLICATION command failed as there was an error when initializing the group communication layer.

- Error: 3097 SQLSTATE: HY000 (ER_GROUP_REPLICATION_COMMUNICATION_LAYER_JOIN_ERROR)

  Message: The START GROUP_REPLICATION command failed as there was an error when joining the communication group.

- Error: 3098 SQLSTATE: HY000 (ER_BEFORE_DML_VALIDATION_ERROR)

  Message: The table does not comply with the requirements by an external plugin.

- Error: 3099 SQLSTATE: HY000 (ER_PREVENTS_VARIABLE_WITHOUT_RBR)

  Message: Cannot change the value of variable %s without binary log format as ROW.

  transaction_write_set_extraction option value is set and binlog_format is not ROW.

- Error: 3100 SQLSTATE: HY000 (ER_RUN_HOOK_ERROR)

  Message: Error on observer while running replication hook '%s'.

- Error: 3101 SQLSTATE: 40000 (ER_TRANSACTION_ROLLBACK_DURING_COMMIT)

  Message: Plugin instructed the server to rollback the current transaction.

  When using Group Replication, this means that a transaction failed the group certification process, due to one or more members detecting a potential conflict, and was thus rolled back. See Chapter 21, *Group Replication*.

- Error: 3102 SQLSTATE: HY000 (ER_GENERATED_COLUMN_FUNCTION_IS_NOT_ALLOWED)

  Message: Expression of generated column '%s' contains a disallowed function.

- Error: 3103 SQLSTATE: HY000 (ER_UNSUPPORTED_ALTER_INPLACE_ON_VIRTUAL_COLUMN)

  Message: INPLACE ADD or DROP of virtual columns cannot be combined with other ALTER TABLE actions

- Error: 3104 SQLSTATE: HY000 (ER_WRONG_FK_OPTION_FOR_GENERATED_COLUMN)

  Message: Cannot define foreign key with %s clause on a generated column.

- Error: 3105 SQLSTATE: HY000 (ER_NON_DEFAULT_VALUE_FOR_GENERATED_COLUMN)

  Message: The value specified for generated column '%s' in table '%s' is not allowed.

- Error: 3106 SQLSTATE: HY000 (ER_UNSUPPORTED_ACTION_ON_GENERATED_COLUMN)

  Message: '%s' is not supported for generated columns.

- Error: 3107 SQLSTATE: HY000 (ER_GENERATED_COLUMN_NON_PRIOR)

  Message: Generated column can refer only to generated columns defined prior to it.

  To address this issue, change the table definition to define each generated column later than any generated columns to which it refers.

- Error: 3108 SQLSTATE: HY000 (ER_DEPENDENT_BY_GENERATED_COLUMN)

  Message: Column '%s' has a generated column dependency.

  You cannot drop or rename a generated column if another column refers to it. You must either drop those columns as well, or redefine them not to refer to the generated column.

- Error: 3109 SQLSTATE: HY000 (ER_GENERATED_COLUMN_REF_AUTO_INC)

Message: Generated column '%s' cannot refer to auto-increment column.

- Error: 3110 SQLSTATE: HY000 (ER_FEATURE_NOT_AVAILABLE)

  Message: The '%s' feature is not available; you need to remove '%s' or use MySQL built with '%s'

- Error: 3111 SQLSTATE: HY000 (ER_CANT_SET_GTID_MODE)

  Message: SET @@GLOBAL.GTID_MODE = %s is not allowed because %s.

- Error: 3112 SQLSTATE: HY000 (ER_CANT_USE_AUTO_POSITION_WITH_GTID_MODE_OFF)

  Message: The replication receiver thread%s cannot start in AUTO_POSITION mode: this server uses @@GLOBAL.GTID_MODE = OFF.

- Error: 3113 SQLSTATE: HY000 (ER_CANT_REPLICATE_ANONYMOUS_WITH_AUTO_POSITION)

  Message: Cannot replicate anonymous transaction when AUTO_POSITION = 1, at file %s, position %lld.

- Error: 3114 SQLSTATE: HY000 (ER_CANT_REPLICATE_ANONYMOUS_WITH_GTID_MODE_ON)

  Message: Cannot replicate anonymous transaction when @@GLOBAL.GTID_MODE = ON, at file %s, position %lld.

- Error: 3115 SQLSTATE: HY000 (ER_CANT_REPLICATE_GTID_WITH_GTID_MODE_OFF)

  Message: Cannot replicate GTID-transaction when @@GLOBAL.GTID_MODE = OFF, at file %s, position %lld.

- Error: 3116 SQLSTATE: HY000 (ER_CANT_ENFORCE_GTID_CONSISTENCY_WITH_ONGOING_GTID_VIOLATING_TX)

  Message: Cannot set ENFORCE_GTID_CONSISTENCY = ON because there are ongoing transactions that violate GTID consistency.

  ER_CANT_SET_ENFORCE_GTID_CONSISTENCY_ON_WITH_ONGOING_GTID_VIOLATING_TRANSACTIONS was renamed to ER_CANT_ENFORCE_GTID_CONSISTENCY_WITH_ONGOING_GTID_VIOLATING_TX.

- Error: 3117 SQLSTATE: HY000 (ER_ENFORCE_GTID_CONSISTENCY_WARN_WITH_ONGOING_GTID_VIOLATING_TX)

  Message: There are ongoing transactions that violate GTID consistency.

  ER_SET_ENFORCE_GTID_CONSISTENCY_WARN_WITH_ONGOING_GTID_VIOLATING_TRANSACTIONS was renamed to ER_ENFORCE_GTID_CONSISTENCY_WARN_WITH_ONGOING_GTID_VIOLATING_TX.

- Error: 3118 SQLSTATE: HY000 (ER_ACCOUNT_HAS_BEEN_LOCKED)

  Message: Access denied for user '%s'@'%s'. Account is locked.

  The account was locked with CREATE USER ... ACCOUNT LOCK or ALTER USER ... ACCOUNT LOCK. An administrator can unlock it with ALTER USER ... ACCOUNT UNLOCK.

- Error: 3119 SQLSTATE: 42000 (ER_WRONG_TABLESPACE_NAME)

  Message: Incorrect tablespace name `%s`

- Error: 3120 SQLSTATE: HY000 (ER_TABLESPACE_IS_NOT_EMPTY)

  Message: Tablespace `%s` is not empty.

- Error: 3121 SQLSTATE: HY000 (ER_WRONG_FILE_NAME)

  Message: Incorrect File Name '%s'.

- Error: 3122 SQLSTATE: HY000 (ER_BOOST_GEOMETRY_INCONSISTENT_TURNS_EXCEPTION)

  Message: Inconsistent intersection points.

- Error: 3123 SQLSTATE: HY000 (ER_WARN_OPTIMIZER_HINT_SYNTAX_ERROR)

  Message: Optimizer hint syntax error

- Error: 3124 SQLSTATE: HY000 (ER_WARN_BAD_MAX_EXECUTION_TIME)

  Message: Unsupported MAX_EXECUTION_TIME

- Error: 3125 SQLSTATE: HY000 (ER_WARN_UNSUPPORTED_MAX_EXECUTION_TIME)

  Message: MAX_EXECUTION_TIME hint is supported by top-level standalone SELECT statements only

  The MAX_EXECUTION_TIME optimizer hint is supported only for SELECT statements.

- Error: 3126 SQLSTATE: HY000 (ER_WARN_CONFLICTING_HINT)

  Message: Hint %s is ignored as conflicting/duplicated

- Error: 3127 SQLSTATE: HY000 (ER_WARN_UNKNOWN_QB_NAME)

  Message: Query block name %s is not found for %s hint

- Error: 3128 SQLSTATE: HY000 (ER_UNRESOLVED_HINT_NAME)

  Message: Unresolved name %s for %s hint

- Error: 3129 SQLSTATE: HY000 (ER_WARN_ON_MODIFYING_GTID_EXECUTED_TABLE)

  Message: Please do not modify the %s table. This is a mysql internal system table to store GTIDs for committed transactions. Modifying it can lead to an inconsistent GTID state.

- Error: 3130 SQLSTATE: HY000 (ER_PLUGGABLE_PROTOCOL_COMMAND_NOT_SUPPORTED)

  Message: Command not supported by pluggable protocols

- Error: 3131 SQLSTATE: 42000 (ER_LOCKING_SERVICE_WRONG_NAME)

  Message: Incorrect locking service lock name '%s'.

  A locking service name was specified as NULL, the empty string, or a string longer than 64 characters. Namespace and lock names must be non-NULL, nonempty, and no more than 64 characters long.

- Error: 3132 SQLSTATE: HY000 (ER_LOCKING_SERVICE_DEADLOCK)

  Message: Deadlock found when trying to get locking service lock; try releasing locks and restarting lock acquisition.

- Error: 3133 SQLSTATE: HY000 (ER_LOCKING_SERVICE_TIMEOUT)

  Message: Service lock wait timeout exceeded.

- Error: 3134 SQLSTATE: HY000 (ER_GIS_MAX_POINTS_IN_GEOMETRY_OVERFLOWED)

  Message: Parameter %s exceeds the maximum number of points in a geometry (%lu) in function %s.

- Error: 3135 SQLSTATE: HY000 (ER_SQL_MODE_MERGED)

  Message: 'NO_ZERO_DATE', 'NO_ZERO_IN_DATE' and 'ERROR_FOR_DIVISION_BY_ZERO' sql modes should be used with strict mode. They will be merged with strict mode in a future release.

- Error: 3136 SQLSTATE: HY000 (ER_VTOKEN_PLUGIN_TOKEN_MISMATCH)

  Message: Version token mismatch for %.*s. Correct value %.*s

  The client has set its version_tokens_session system variable to the list of tokens it requires the server to match, but the server token list has at least one matching token name that has a value different from what the client requires. See Section 5.6.5, “Version Tokens”.

- Error: 3137 SQLSTATE: HY000 (ER_VTOKEN_PLUGIN_TOKEN_NOT_FOUND)

  Message: Version token %.*s not found.

  The client has set its version_tokens_session system variable to the list of tokens it requires the server to match, but the server token list is missing at least one of those tokens. See Section 5.6.5, “Version Tokens”.

- Error: 3138 SQLSTATE: HY000 (ER_CANT_SET_VARIABLE_WHEN_OWNING_GTID)

  Message: Variable %s cannot be changed by a client that owns a GTID. The client owns %s. Ownership is released on COMMIT or ROLLBACK.

- Error: 3139 SQLSTATE: HY000 (ER_SLAVE_CHANNEL_OPERATION_NOT_ALLOWED)

  Message: %s cannot be performed on channel '%s'.

- Error: 3140 SQLSTATE: 22032 (ER_INVALID_JSON_TEXT)

  Message: Invalid JSON text: "%s" at position %u in value for column '%s'.

- Error: 3141 SQLSTATE: 22032 (ER_INVALID_JSON_TEXT_IN_PARAM)

  Message: Invalid JSON text in argument %u to function %s: "%s" at position %u.%s

- Error: 3142 SQLSTATE: HY000 (ER_INVALID_JSON_BINARY_DATA)

  Message: The JSON binary value contains invalid data.

- Error: 3143 SQLSTATE: 42000 (ER_INVALID_JSON_PATH)

  Message: Invalid JSON path expression. The error is around character position %u.%s

- Error: 3144 SQLSTATE: 22032 (ER_INVALID_JSON_CHARSET)

  Message: Cannot create a JSON value from a string with CHARACTER SET '%s'.

- Error: 3145 SQLSTATE: 22032 (ER_INVALID_JSON_CHARSET_IN_FUNCTION)

  Message: Invalid JSON character data provided to function %s: '%s'; utf8 is required.

- Error: 3146 SQLSTATE: 22032 (ER_INVALID_TYPE_FOR_JSON)

  Message: Invalid data type for JSON data in argument %u to function %s; a JSON string or JSON type is required.

- Error: 3147 SQLSTATE: 22032 (ER_INVALID_CAST_TO_JSON)

  Message: Cannot CAST value to JSON.

- Error: 3148 SQLSTATE: 42000 (ER_INVALID_JSON_PATH_CHARSET)

Message: A path expression must be encoded in the utf8 character set. The path expression '%s' is encoded in character set '%s'.

- Error: 3149 SQLSTATE: 42000 (ER_INVALID_JSON_PATH_WILDCARD)

  Message: In this situation, path expressions may not contain the * and ** tokens or an array range.

- Error: 3150 SQLSTATE: 22032 (ER_JSON_VALUE_TOO_BIG)

  Message: The JSON value is too big to be stored in a JSON column.

- Error: 3151 SQLSTATE: 22032 (ER_JSON_KEY_TOO_BIG)

  Message: The JSON object contains a key name that is too long.

- Error: 3152 SQLSTATE: 42000 (ER_JSON_USED_AS_KEY)

  Message: JSON column '%s' supports indexing only via generated columns on a specified JSON path.

- Error: 3153 SQLSTATE: 42000 (ER_JSON_VACUOUS_PATH)

  Message: The path expression '$' is not allowed in this context.

- Error: 3154 SQLSTATE: 42000 (ER_JSON_BAD_ONE_OR_ALL_ARG)

  Message: The oneOrAll argument to %s may take these values: 'one' or 'all'.

- Error: 3155 SQLSTATE: 22003 (ER_NUMERIC_JSON_VALUE_OUT_OF_RANGE)

  Message: Out of range JSON value for CAST to %s%s from column %s at row %ld

- Error: 3156 SQLSTATE: 22018 (ER_INVALID_JSON_VALUE_FOR_CAST)

  Message: Invalid JSON value for CAST to %s%s from column %s at row %ld

- Error: 3157 SQLSTATE: 22032 (ER_JSON_DOCUMENT_TOO_DEEP)

  Message: The JSON document exceeds the maximum depth.

- Error: 3158 SQLSTATE: 22032 (ER_JSON_DOCUMENT_NULL_KEY)

  Message: JSON documents may not contain NULL member names.

- Error: 3159 SQLSTATE: HY000 (ER_SECURE_TRANSPORT_REQUIRED)

  Message: Connections using insecure transport are prohibited while --require_secure_transport=ON.

  With the require_secure_transport system variable, clients can connect only using secure transports. Qualifying connections are those using SSL, a Unix socket file, or shared memory.

- Error: 3160 SQLSTATE: HY000 (ER_NO_SECURE_TRANSPORTS_CONFIGURED)

  Message: No secure transports (SSL or Shared Memory) are configured, unable to set --require_secure_transport=ON.

  The require_secure_transport system variable cannot be enabled if the server does not support at least one secure transport. Configure the server with the required SSL keys/certificates to enable SSL connections, or enable the shared_memory system variable to enable shared-memory connections.

- Error: 3161 SQLSTATE: HY000 (ER_DISABLED_STORAGE_ENGINE)

Message: Storage engine %s is disabled (Table creation is disallowed).

An attempt was made to create a table or tablespace using a storage engine listed in the value of the `disabled_storage_engines` system variable, or to change an existing table or tablespace to such an engine. Choose a different storage engine.

- Error: `3162` SQLSTATE: `HY000` (`ER_USER_DOES_NOT_EXIST`)

  Message: Authorization ID %s does not exist.

- Error: `3163` SQLSTATE: `HY000` (`ER_USER_ALREADY_EXISTS`)

  Message: Authorization ID %s already exists.

- Error: `3164` SQLSTATE: `HY000` (`ER_AUDIT_API_ABORT`)

  Message: Aborted by Audit API ('%s';%d).

  This error indicates that an audit plugin terminated execution of an event. The message typically indicates the event subclass name and a numeric status value.

- Error: `3165` SQLSTATE: `42000` (`ER_INVALID_JSON_PATH_ARRAY_CELL`)

  Message: A path expression is not a path to a cell in an array.

- Error: `3166` SQLSTATE: `HY000` (`ER_BUFPOOL_RESIZE_INPROGRESS`)

  Message: Another buffer pool resize is already in progress.

- Error: `3167` SQLSTATE: `HY000` (`ER_FEATURE_DISABLED_SEE_DOC`)

  Message: The '%s' feature is disabled; see the documentation for '%s'

- Error: `3168` SQLSTATE: `HY000` (`ER_SERVER_ISNT_AVAILABLE`)

  Message: Server isn't available

- Error: `3169` SQLSTATE: `HY000` (`ER_SESSION_WAS_KILLED`)

  Message: Session was killed

- Error: `3170` SQLSTATE: `HY000` (`ER_CAPACITY_EXCEEDED`)

  Message: Memory capacity of %llu bytes for '%s' exceeded. %s

- Error: `3171` SQLSTATE: `HY000` (`ER_CAPACITY_EXCEEDED_IN_RANGE_OPTIMIZER`)

  Message: Range optimization was not done for this query.

- Error: `3173` SQLSTATE: `HY000` (`ER_CANT_WAIT_FOR_EXECUTED_GTID_SET_WHILE_OWNING_A_GTID`)

  Message: The client holds ownership of the GTID %s. Therefore, WAIT_FOR_EXECUTED_GTID_SET cannot wait for this GTID.

- Error: `3174` SQLSTATE: `HY000` (`ER_CANNOT_ADD_FOREIGN_BASE_COL_VIRTUAL`)

  Message: Cannot add foreign key on the base column of indexed virtual column.

- Error: `3175` SQLSTATE: `HY000` (`ER_CANNOT_CREATE_VIRTUAL_INDEX_CONSTRAINT`)

  Message: Cannot create index on virtual column whose base column has foreign constraint.

- Error: 3176 SQLSTATE: HY000 (ER_ERROR_ON_MODIFYING_GTID_EXECUTED_TABLE)

  Message: Please do not modify the %s table with an XA transaction. This is an internal system table used to store GTIDs for committed transactions. Although modifying it can lead to an inconsistent GTID state, if neccessary you can modify it with a non-XA transaction.

- Error: 3177 SQLSTATE: HY000 (ER_LOCK_REFUSED_BY_ENGINE)

  Message: Lock acquisition refused by storage engine.

- Error: 3178 SQLSTATE: HY000 (ER_UNSUPPORTED_ALTER_ONLINE_ON_VIRTUAL_COLUMN)

  Message: ADD COLUMN col...VIRTUAL, ADD INDEX(col)

- Error: 3179 SQLSTATE: HY000 (ER_MASTER_KEY_ROTATION_NOT_SUPPORTED_BY_SE)

  Message: Master key rotation is not supported by storage engine.

- Error: 3181 SQLSTATE: HY000 (ER_MASTER_KEY_ROTATION_BINLOG_FAILED)

  Message: Write to binlog failed. However, master key rotation has been completed successfully.

- Error: 3182 SQLSTATE: HY000 (ER_MASTER_KEY_ROTATION_SE_UNAVAILABLE)

  Message: Storage engine is not available.

- Error: 3183 SQLSTATE: HY000 (ER_TABLESPACE_CANNOT_ENCRYPT)

  Message: This tablespace can't be encrypted.

- Error: 3184 SQLSTATE: HY000 (ER_INVALID_ENCRYPTION_OPTION)

  Message: Invalid encryption option.

- Error: 3185 SQLSTATE: HY000 (ER_CANNOT_FIND_KEY_IN_KEYRING)

  Message: Can't find master key from keyring, please check keyring plugin is loaded.

- Error: 3186 SQLSTATE: HY000 (ER_CAPACITY_EXCEEDED_IN_PARSER)

  Message: Parser bailed out for this query.

- Error: 3187 SQLSTATE: HY000 (ER_UNSUPPORTED_ALTER_ENCRYPTION_INPLACE)

  Message: Cannot alter encryption attribute by inplace algorithm.

- Error: 3188 SQLSTATE: HY000 (ER_KEYRING_UDF_KEYRING_SERVICE_ERROR)

  Message: Function '%s' failed because underlying keyring service returned an error. Please check if a keyring plugin is installed and that provided arguments are valid for the keyring you are using.

- Error: 3189 SQLSTATE: HY000 (ER_USER_COLUMN_OLD_LENGTH)

  Message: It seems that your db schema is old. The %s column is 77 characters long and should be 93 characters long. Please run mysql_upgrade.

- Error: 3190 SQLSTATE: HY000 (ER_CANT_RESET_MASTER)

  Message: RESET MASTER is not allowed because %s.

- Error: 3191 SQLSTATE: HY000 (ER_GROUP_REPLICATION_MAX_GROUP_SIZE)

  Message: The START GROUP_REPLICATION command failed since the group already has 9 members.

- Error: 3192 SQLSTATE: HY000 (ER_CANNOT_ADD_FOREIGN_BASE_COL_STORED)

  Message: Cannot add foreign key on the base column of stored column.

- Error: 3193 SQLSTATE: HY000 (ER_TABLE_REFERENCED)

  Message: Cannot complete the operation because table is referenced by another connection.

- Error: 3197 SQLSTATE: HY000 (ER_XA_RETRY)

  Message: The resource manager is not able to commit the transaction branch at this time. Please retry later.

  ER_XA_RETRY was added in 8.0.2.

- Error: 3198 SQLSTATE: HY000 (ER_KEYRING_AWS_UDF_AWS_KMS_ERROR)

  Message: Function %s failed due to: %s.

  ER_KEYRING_AWS_UDF_AWS_KMS_ERROR was added in 8.0.2.

- Error: 3199 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_XA)

  Message: Statement is unsafe because it is being used inside a XA transaction. Concurrent XA transactions may deadlock on slaves when replicated using statements.

  ER_BINLOG_UNSAFE_XA was added in 8.0.2.

- Error: 3500 SQLSTATE: HY000 (ER_UNSUPPORT_COMPRESSED_TEMPORARY_TABLE)

  Message: CREATE TEMPORARY TABLE is not allowed with ROW_FORMAT=COMPRESSED or KEY_BLOCK_SIZE.

- Error: 3501 SQLSTATE: HY000 (ER_ACL_OPERATION_FAILED)

  Message: The ACL operation failed due to the following error from SE: errcode %d - %s

- Error: 3502 SQLSTATE: HY000 (ER_UNSUPPORTED_INDEX_ALGORITHM)

  Message: This storage engine does not support the %s index algorithm, storage engine default was used instead.

- Error: 3503 SQLSTATE: 42Y07 (ER_NO_SUCH_DB)

  Message: Database '%s' doesn't exist

- Error: 3504 SQLSTATE: HY000 (ER_TOO_BIG_ENUM)

  Message: Too many enumeration values for column %s.

- Error: 3505 SQLSTATE: HY000 (ER_TOO_LONG_SET_ENUM_VALUE)

  Message: Too long enumeration/set value for column %s.

- Error: 3506 SQLSTATE: HY000 (ER_INVALID_DD_OBJECT)

  Message: %s dictionary object is invalid. (%s)

- Error: 3507 SQLSTATE: HY000 (ER_UPDATING_DD_TABLE)

  Message: Failed to update %s dictionary object.

- Error: 3508 SQLSTATE: HY000 (ER_INVALID_DD_OBJECT_ID)

  Message: Dictionary object id (%lu) does not exist.

- Error: 3509 SQLSTATE: HY000 (ER_INVALID_DD_OBJECT_NAME)

  Message: Dictionary object name '%s' is invalid. (%s)

- Error: 3510 SQLSTATE: HY000 (ER_TABLESPACE_MISSING_WITH_NAME)

  Message: Tablespace %s doesn't exist.

- Error: 3511 SQLSTATE: HY000 (ER_TOO_LONG_ROUTINE_COMMENT)

  Message: Comment for routine '%s' is too long (max = %lu)

- Error: 3512 SQLSTATE: HY000 (ER_SP_LOAD_FAILED)

  Message: Failed to load routine '%s'.

- Error: 3513 SQLSTATE: HY000 (ER_INVALID_BITWISE_OPERANDS_SIZE)

  Message: Binary operands of bitwise operators must be of equal length

- Error: 3514 SQLSTATE: HY000 (ER_INVALID_BITWISE_AGGREGATE_OPERANDS_SIZE)

  Message: Aggregate bitwise functions cannot accept arguments longer than 511 bytes; consider using the SUBSTRING() function

- Error: 3515 SQLSTATE: HY000 (ER_WARN_UNSUPPORTED_HINT)

  Message: Hints aren't supported in %s

- Error: 3516 SQLSTATE: 22S01 (ER_UNEXPECTED_GEOMETRY_TYPE)

  Message: %s value is a geometry of unexpected type %s in %s.

- Error: 3517 SQLSTATE: SR002 (ER_SRS_PARSE_ERROR)

  Message: Can't parse the spatial reference system definition of SRID %u.

- Error: 3518 SQLSTATE: SR003 (ER_SRS_PROJ_PARAMETER_MISSING)

  Message: The spatial reference system definition for SRID %u does not specify the mandatory %s (EPSG %u) projection parameter.

- Error: 3519 SQLSTATE: 01000 (ER_WARN_SRS_NOT_FOUND)

  Message: There's no spatial reference system with SRID %u. Calculations in function %s will be done in SRID 0.

- Error: 3520 SQLSTATE: 22S00 (ER_SRS_NOT_CARTESIAN)

  Message: Function %s is only defined for Cartesian spatial reference systems, but one of its arguments is in SRID %u, which is not Cartesian.

- Error: 3521 SQLSTATE: SR001 (ER_SRS_NOT_CARTESIAN_UNDEFINED)

  Message: Function %s is only defined for Cartesian spatial reference systems, but one of its arguments is in SRID %u, which has not been defined.

- Error: 3522 SQLSTATE: HY000 (ER_PK_INDEX_CANT_BE_INVISIBLE)

  Message: A primary key index cannot be invisible

- Error: 3523 SQLSTATE: HY000 (ER_UNKNOWN_AUTHID)

  Message: Unknown authorization ID `%s`@`%s`

- Error: 3524 SQLSTATE: HY000 (ER_FAILED_ROLE_GRANT)

  Message: Failed to grant %s` to %s

- Error: 3525 SQLSTATE: HY000 (ER_OPEN_ROLE_TABLES)

  Message: Failed to open the security system tables

- Error: 3526 SQLSTATE: HY000 (ER_FAILED_DEFAULT_ROLES)

  Message: Failed to set default roles

- Error: 3527 SQLSTATE: HY000 (ER_COMPONENTS_NO_SCHEME)

  Message: Cannot find schema in specified URN: '%s'.

- Error: 3528 SQLSTATE: HY000 (ER_COMPONENTS_NO_SCHEME_SERVICE)

  Message: Cannot acquire scheme load service implementation for schema '%s' in specified URN: '%s'.

- Error: 3529 SQLSTATE: HY000 (ER_COMPONENTS_CANT_LOAD)

  Message: Cannot load component from specified URN: '%s'.

- Error: 3530 SQLSTATE: HY000 (ER_ROLE_NOT_GRANTED)

  Message: `%s`@`%s` is not granted to `%s`@`%s`

- Error: 3531 SQLSTATE: HY000 (ER_FAILED_REVOKE_ROLE)

  Message: Could not revoke role from `%s`@`%s`

- Error: 3532 SQLSTATE: HY000 (ER_RENAME_ROLE)

  Message: Renaming of a role identifier is forbidden

- Error: 3533 SQLSTATE: HY000 (ER_COMPONENTS_CANT_ACQUIRE_SERVICE_IMPLEMENTATION)

  Message: Cannot acquire specified service implementation: '%s'.

- Error: 3534 SQLSTATE: HY000 (ER_COMPONENTS_CANT_SATISFY_DEPENDENCY)

  Message: Cannot satisfy dependency for service '%s' required by component '%s'.

- Error: 3535 SQLSTATE: HY000 (ER_COMPONENTS_LOAD_CANT_REGISTER_SERVICE_IMPLEMENTATION)

  Message: Cannot register service implementation '%s' provided by component '%s'.

- Error: 3536 SQLSTATE: HY000 (ER_COMPONENTS_LOAD_CANT_INITIALIZE)

  Message: Initialization method provided by component '%s' failed.

- Error: 3537 SQLSTATE: HY000 (ER_COMPONENTS_UNLOAD_NOT_LOADED)

  Message: Component specified by URN '%s' to unload has not been loaded before.

- Error: 3538 SQLSTATE: HY000 (ER_COMPONENTS_UNLOAD_CANT_DEINITIALIZE)

  Message: De-initialization method provided by component '%s' failed.

- Error: 3539 SQLSTATE: HY000 (ER_COMPONENTS_CANT_RELEASE_SERVICE)

Message: Release of previously acquired service implementation failed.

- Error: 3540 SQLSTATE: HY000 (ER_COMPONENTS_UNLOAD_CANT_UNREGISTER_SERVICE)

  Message: Unregistration of service implementation '%s' provided by component '%s' failed during unloading of the component.

- Error: 3541 SQLSTATE: HY000 (ER_COMPONENTS_CANT_UNLOAD)

  Message: Cannot unload component from specified URN: '%s'.

- Error: 3542 SQLSTATE: HY000 (ER_WARN_UNLOAD_THE_NOT_PERSISTED)

  Message: The Persistent Dynamic Loader was used to unload a component '%s', but it was not used to load that component before.

- Error: 3543 SQLSTATE: HY000 (ER_COMPONENT_TABLE_INCORRECT)

  Message: The mysql.component table is missing or has an incorrect definition.

- Error: 3544 SQLSTATE: HY000 (ER_COMPONENT_MANIPULATE_ROW_FAILED)

  Message: Failed to manipulate component '%s' persistence data. Error code %d from storage engine.

- Error: 3545 SQLSTATE: HY000 (ER_COMPONENTS_UNLOAD_DUPLICATE_IN_GROUP)

  Message: The component with specified URN: '%s' was specified in group more than once.

- Error: 3546 SQLSTATE: HY000 (ER_CANT_SET_GTID_PURGED_DUE_SETS_CONSTRAINTS)

  Message: @@GLOBAL.GTID_PURGED cannot be changed: %s

- Error: 3547 SQLSTATE: HY000 (ER_CANNOT_LOCK_USER_MANAGEMENT_CACHES)

  Message: Can not lock user management caches for processing.

- Error: 3548 SQLSTATE: SR001 (ER_SRS_NOT_FOUND)

  Message: There's no spatial reference system with SRID %u.

- Error: 3549 SQLSTATE: HY000 (ER_VARIABLE_NOT_PERSISTED)

  Message: Variables cannot be persisted. Please retry.

- Error: 3550 SQLSTATE: HY000 (ER_IS_QUERY_INVALID_CLAUSE)

  Message: Information schema queries do not support the '%s' clause.

- Error: 3551 SQLSTATE: HY000 (ER_UNABLE_TO_STORE_STATISTICS)

  Message: Unable to store dynamic %s statistics into data dictionary.

- Error: 3552 SQLSTATE: HY000 (ER_NO_SYSTEM_SCHEMA_ACCESS)

  Message: Access to system schema '%s' is rejected.

- Error: 3553 SQLSTATE: HY000 (ER_NO_SYSTEM_TABLESPACE_ACCESS)

  Message: Access to system tablespace '%s' is rejected.

- Error: 3554 SQLSTATE: HY000 (ER_NO_SYSTEM_TABLE_ACCESS)

  Message: Access to %s '%s.%s' is rejected.

- Error: 3555 SQLSTATE: HY000 (ER_NO_SYSTEM_TABLE_ACCESS_FOR_DICTIONARY_TABLE)

  Message: data dictionary table

  ER_NO_SYSTEM_TABLE_ACCESS_FOR_DICTIONARY_TABLE was added in 8.0.1.
- Error: 3556 SQLSTATE: HY000 (ER_NO_SYSTEM_TABLE_ACCESS_FOR_SYSTEM_TABLE)

  Message: system table

  ER_NO_SYSTEM_TABLE_ACCESS_FOR_SYSTEM_TABLE was added in 8.0.1.
- Error: 3557 SQLSTATE: HY000 (ER_NO_SYSTEM_TABLE_ACCESS_FOR_TABLE)

  Message: table

  ER_NO_SYSTEM_TABLE_ACCESS_FOR_TABLE was added in 8.0.1.
- Error: 3558 SQLSTATE: 22023 (ER_INVALID_OPTION_KEY)

  Message: Invalid option key '%s' in function %s.

  ER_INVALID_OPTION_KEY was added in 8.0.1.
- Error: 3559 SQLSTATE: 22023 (ER_INVALID_OPTION_VALUE)

  Message: Invalid value '%s' for option '%s' in function '%s'.

  ER_INVALID_OPTION_VALUE was added in 8.0.1.
- Error: 3560 SQLSTATE: 22023 (ER_INVALID_OPTION_KEY_VALUE_PAIR)

  Message: The string '%s' is not a valid key %c value pair in function %s.

  ER_INVALID_OPTION_KEY_VALUE_PAIR was added in 8.0.1.
- Error: 3561 SQLSTATE: 22023 (ER_INVALID_OPTION_START_CHARACTER)

  Message: The options argument in function %s starts with the invalid character '%c'.

  ER_INVALID_OPTION_START_CHARACTER was added in 8.0.1.
- Error: 3562 SQLSTATE: 22023 (ER_INVALID_OPTION_END_CHARACTER)

  Message: The options argument in function %s ends with the invalid character '%c'.

  ER_INVALID_OPTION_END_CHARACTER was added in 8.0.1.
- Error: 3563 SQLSTATE: 22023 (ER_INVALID_OPTION_CHARACTERS)

  Message: The options argument in function %s contains the invalid character sequence '%s'.

  ER_INVALID_OPTION_CHARACTERS was added in 8.0.1.
- Error: 3564 SQLSTATE: 22023 (ER_DUPLICATE_OPTION_KEY)

  Message: Duplicate option key '%s' in funtion '%s'.

  ER_DUPLICATE_OPTION_KEY was added in 8.0.1.
- Error: 3565 SQLSTATE: 01000 (ER_WARN_SRS_NOT_FOUND_AXIS_ORDER)

  Message: There's no spatial reference system with SRID %u. The axis order is unknown.

  ER_WARN_SRS_NOT_FOUND_AXIS_ORDER was added in 8.0.1.

- Error: 3566 SQLSTATE: HY000 (ER_NO_ACCESS_TO_NATIVE_FCT)

  Message: Access to native function '%s' is rejected.

  ER_NO_ACCESS_TO_NATIVE_FCT was added in 8.0.1.

- Error: 3567 SQLSTATE: HY000 (ER_RESET_MASTER_TO_VALUE_OUT_OF_RANGE)

  Message: The requested value '%lu' for the next binary log index is out of range. Please use a value between '1' and '%lu'.

  ER_RESET_MASTER_TO_VALUE_OUT_OF_RANGE was added in 8.0.1.

- Error: 3568 SQLSTATE: HY000 (ER_UNRESOLVED_TABLE_LOCK)

  Message: Unresolved table name %s in locking clause.

  ER_UNRESOLVED_TABLE_LOCK was added in 8.0.1.

- Error: 3569 SQLSTATE: HY000 (ER_DUPLICATE_TABLE_LOCK)

  Message: Table %s appears in multiple locking clauses.

  ER_DUPLICATE_TABLE_LOCK was added in 8.0.1.

- Error: 3570 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_SKIP_LOCKED)

  Message: Statement is unsafe because it uses SKIP LOCKED. The set of inserted values is non-deterministic.

  ER_BINLOG_UNSAFE_SKIP_LOCKED was added in 8.0.1.

- Error: 3571 SQLSTATE: HY000 (ER_BINLOG_UNSAFE_NOWAIT)

  Message: Statement is unsafe because it uses NOWAIT. Whether the command will succeed or fail is not deterministic.

  ER_BINLOG_UNSAFE_NOWAIT was added in 8.0.1.

- Error: 3572 SQLSTATE: HY000 (ER_LOCK_NOWAIT)

  Message: Statement aborted because lock(s) could not be acquired immediately and NOWAIT is set.

  ER_LOCK_NOWAIT was added in 8.0.1.

- Error: 3573 SQLSTATE: HY000 (ER_CTE_RECURSIVE_REQUIRES_UNION)

  Message: Recursive Common Table Expression '%s' should contain a UNION

  ER_CTE_RECURSIVE_REQUIRES_UNION was added in 8.0.1.

- Error: 3574 SQLSTATE: HY000 (ER_CTE_RECURSIVE_REQUIRES_NONRECURSIVE_FIRST)

  Message: Recursive Common Table Expression '%s' should have one or more non-recursive query blocks followed by one or more recursive ones

  ER_CTE_RECURSIVE_REQUIRES_NONRECURSIVE_FIRST was added in 8.0.1.

- Error: 3575 SQLSTATE: HY000 (ER_CTE_RECURSIVE_FORBIDS_AGGREGATION)

  Message: Recursive Common Table Expression '%s' can contain neither aggregation nor window functions in recursive query block

  ER_CTE_RECURSIVE_FORBIDS_AGGREGATION was added in 8.0.1.

- Error: 3576 SQLSTATE: HY000 (ER_CTE_RECURSIVE_FORBIDDEN_JOIN_ORDER)

  Message: In recursive query block of Recursive Common Table Expression '%s', the recursive table must neither be in the right argument of a LEFT JOIN, nor be forced to be non-first with join order hints

  ER_CTE_RECURSIVE_FORBIDDEN_JOIN_ORDER was added in 8.0.1.

- Error: 3577 SQLSTATE: HY000 (ER_CTE_RECURSIVE_REQUIRES_SINGLE_REFERENCE)

  Message: In recursive query block of Recursive Common Table Expression '%s', the recursive table must be referenced only once, and not in any subquery

  ER_CTE_RECURSIVE_REQUIRES_SINGLE_REFERENCE was added in 8.0.1.

- Error: 3578 SQLSTATE: HY000 (ER_SWITCH_TMP_ENGINE)

  Message: '%s' requires @@internal_tmp_disk_storage_engine=InnoDB

  ER_SWITCH_TMP_ENGINE was added in 8.0.1.

- Error: 3579 SQLSTATE: HY000 (ER_WINDOW_NO_SUCH_WINDOW)

  Message: Window name '%s' is not defined.

  ER_WINDOW_NO_SUCH_WINDOW was added in 8.0.2.

- Error: 3580 SQLSTATE: HY000 (ER_WINDOW_CIRCULARITY_IN_WINDOW_GRAPH)

  Message: There is a circularity in the window dependency graph.

  ER_WINDOW_CIRCULARITY_IN_WINDOW_GRAPH was added in 8.0.2.

- Error: 3581 SQLSTATE: HY000 (ER_WINDOW_NO_CHILD_PARTITIONING)

  Message: A window which depends on another cannot define partitioning.

  ER_WINDOW_NO_CHILD_PARTITIONING was added in 8.0.2.

- Error: 3582 SQLSTATE: HY000 (ER_WINDOW_NO_INHERIT_FRAME)

  Message: Window '%s' has a frame definition, so cannot be referenced by another window.

  ER_WINDOW_NO_INHERIT_FRAME was added in 8.0.2.

- Error: 3583 SQLSTATE: HY000 (ER_WINDOW_NO_REDEFINE_ORDER_BY)

  Message: Window '%s' cannot inherit '%s' since both contain an ORDER BY clause.

  ER_WINDOW_NO_REDEFINE_ORDER_BY was added in 8.0.2.

- Error: 3584 SQLSTATE: HY000 (ER_WINDOW_FRAME_START_ILLEGAL)

  Message: Window '%s': frame start cannot be UNBOUNDED FOLLOWING.

  ER_WINDOW_FRAME_START_ILLEGAL was added in 8.0.2.

- Error: 3585 SQLSTATE: HY000 (ER_WINDOW_FRAME_END_ILLEGAL)

  Message: Window '%s': frame end cannot be UNBOUNDED PRECEDING.

  ER_WINDOW_FRAME_END_ILLEGAL was added in 8.0.2.

- Error: 3586 SQLSTATE: HY000 (ER_WINDOW_FRAME_ILLEGAL)

Message: Window '%s': frame start or end is negative or NULL

ER_WINDOW_FRAME_ILLEGAL was added in 8.0.2.

- Error: 3587 SQLSTATE: HY000 (ER_WINDOW_RANGE_FRAME_ORDER_TYPE)

  Message: Window '%s' with RANGE N PRECEDING/FOLLOWING frame requires exactly one ORDER BY expression, of numeric or temporal type

  ER_WINDOW_RANGE_FRAME_ORDER_TYPE was added in 8.0.2.

- Error: 3588 SQLSTATE: HY000 (ER_WINDOW_RANGE_FRAME_TEMPORAL_TYPE)

  Message: Window '%s' with RANGE frame has ORDER BY expression of datetime type. Only INTERVAL bound value allowed.

  ER_WINDOW_RANGE_FRAME_TEMPORAL_TYPE was added in 8.0.2.

- Error: 3589 SQLSTATE: HY000 (ER_WINDOW_RANGE_FRAME_NUMERIC_TYPE)

  Message: Window '%s' with RANGE frame has ORDER BY expression of numeric type, INTERVAL bound value not allowed.

  ER_WINDOW_RANGE_FRAME_NUMERIC_TYPE was added in 8.0.2.

- Error: 3590 SQLSTATE: HY000 (ER_WINDOW_RANGE_BOUND_NOT_CONSTANT)

  Message: Window '%s' has a non-constant frame bound.

  ER_WINDOW_RANGE_BOUND_NOT_CONSTANT was added in 8.0.2.

- Error: 3591 SQLSTATE: HY000 (ER_WINDOW_DUPLICATE_NAME)

  Message: Window '%s' is defined twice.

  ER_WINDOW_DUPLICATE_NAME was added in 8.0.2.

- Error: 3592 SQLSTATE: HY000 (ER_WINDOW_ILLEGAL_ORDER_BY)

  Message: Window '%s': ORDER BY or PARTITION BY uses legacy position indication which is not supported, use expression.

  ER_WINDOW_ILLEGAL_ORDER_BY was added in 8.0.2.

- Error: 3593 SQLSTATE: HY000 (ER_WINDOW_INVALID_WINDOW_FUNC_USE)

  Message: You cannot use the window function '%s' in this context.'

  ER_WINDOW_INVALID_WINDOW_FUNC_USE was added in 8.0.2.

- Error: 3594 SQLSTATE: HY000 (ER_WINDOW_INVALID_WINDOW_FUNC_ALIAS_USE)

  Message: You cannot use the alias '%s' of an expression containing a window function in this context.'

  ER_WINDOW_INVALID_WINDOW_FUNC_ALIAS_USE was added in 8.0.2.

- Error: 3595 SQLSTATE: HY000 (ER_WINDOW_NESTED_WINDOW_FUNC_USE_IN_WINDOW_SPEC)

  Message: You cannot nest a window function in the specification of window '%s'.

  ER_WINDOW_NESTED_WINDOW_FUNC_USE_IN_WINDOW_SPEC was added in 8.0.2.

- Error: 3596 SQLSTATE: HY000 (ER_WINDOW_ROWS_INTERVAL_USE)

  Message: Window '%s': INTERVAL can only be used with RANGE frames.

  ER_WINDOW_ROWS_INTERVAL_USE was added in 8.0.2.

- Error: 3597 SQLSTATE: HY000 (ER_WINDOW_NO_GROUP_ORDER)

  Message: ASC or DESC with GROUP BY isn't allowed with window functions; put ASC or DESC in ORDER BY

  ER_WINDOW_NO_GROUP_ORDER was added in 8.0.2.

- Error: 3598 SQLSTATE: HY000 (ER_WINDOW_EXPLAIN_JSON)

  Message: To get information about window functions use EXPLAIN FORMAT=JSON

  ER_WINDOW_EXPLAIN_JSON was added in 8.0.2.

- Error: 3599 SQLSTATE: HY000 (ER_WINDOW_FUNCTION_IGNORES_FRAME)

  Message: Window function '%s' ignores the frame clause of window '%s' and aggregates over the whole partition

  ER_WINDOW_FUNCTION_IGNORES_FRAME was added in 8.0.2.

- Error: 3600 SQLSTATE: HY000 (ER_WINDOW_SE_NOT_ACCEPTABLE)

  Message: Windowing requires @@internal_tmp_mem_storage_engine=TempTable.

  ER_WINDOW_SE_NOT_ACCEPTABLE was added in 8.0.2.

- Error: 3601 SQLSTATE: HY000 (ER_INVALID_NO_OF_ARGS)

  Message: Too many arguments for function %s: %lu; maximum allowed is %s.

  ER_INVALID_NO_OF_ARGS was added in 8.0.1.

- Error: 3602 SQLSTATE: HY000 (ER_FIELD_IN_GROUPING_NOT_GROUP_BY)

  Message: Argument #%u of GROUPING function is not in GROUP BY

  ER_FIELD_IN_GROUPING_NOT_GROUP_BY was added in 8.0.1.

- Error: 3603 SQLSTATE: HY000 (ER_TOO_LONG_TABLESPACE_COMMENT)

  Message: Comment for tablespace '%s' is too long (max = %lu)

  ER_TOO_LONG_TABLESPACE_COMMENT was added in 8.0.1.

- Error: 3604 SQLSTATE: HY000 (ER_ENGINE_CANT_DROP_TABLE)

  Message: Storage engine can't drop table '%s'

  ER_ENGINE_CANT_DROP_TABLE was added in 8.0.1.

- Error: 3605 SQLSTATE: HY000 (ER_ENGINE_CANT_DROP_MISSING_TABLE)

  Message: Storage engine can't drop table '%s' because it is missing. Use DROP TABLE IF EXISTS to remove it from data-dictionary.

  ER_ENGINE_CANT_DROP_MISSING_TABLE was added in 8.0.1.

- Error: 3606 SQLSTATE: HY000 (ER_TABLESPACE_DUP_FILENAME)

Message: Duplicate file name for tablespace '%s'

ER_TABLESPACE_DUP_FILENAME was added in 8.0.1.

- Error: 3607 SQLSTATE: HY000 (ER_DB_DROP_RMDIR2)

  Message: Problem while dropping database. Can't remove database directory (%s). Please remove it manually.

  ER_DB_DROP_RMDIR2 was added in 8.0.1.

- Error: 3608 SQLSTATE: HY000 (ER_IMP_NO_FILES_MATCHED)

  Message: No SDI files matched the pattern '%s'

  ER_IMP_NO_FILES_MATCHED was added in 8.0.1.

- Error: 3609 SQLSTATE: HY000 (ER_IMP_SCHEMA_DOES_NOT_EXIST)

  Message: Schema '%s', referenced in SDI, does not exist.

  ER_IMP_SCHEMA_DOES_NOT_EXIST was added in 8.0.1.

- Error: 3610 SQLSTATE: HY000 (ER_IMP_TABLE_ALREADY_EXISTS)

  Message: Table '%s.%s', referenced in SDI, already exists.

  ER_IMP_TABLE_ALREADY_EXISTS was added in 8.0.1.

- Error: 3611 SQLSTATE: HY000 (ER_IMP_INCOMPATIBLE_MYSQLD_VERSION)

  Message: Imported mysqld_version (%llu) is not compatible with current (%llu)

  ER_IMP_INCOMPATIBLE_MYSQLD_VERSION was added in 8.0.1.

- Error: 3612 SQLSTATE: HY000 (ER_IMP_INCOMPATIBLE_DD_VERSION)

  Message: Imported dd version (%u) is not compatible with current (%u)

  ER_IMP_INCOMPATIBLE_DD_VERSION was added in 8.0.1.

- Error: 3613 SQLSTATE: HY000 (ER_IMP_INCOMPATIBLE_SDI_VERSION)

  Message: Imported sdi version (%llu) is not compatible with current (%llu)

  ER_IMP_INCOMPATIBLE_SDI_VERSION was added in 8.0.1.

- Error: 3614 SQLSTATE: HY000 (ER_WARN_INVALID_HINT)

  Message: Invalid number of arguments for hint %s

  ER_WARN_INVALID_HINT was added in 8.0.1.

- Error: 3615 SQLSTATE: HY000 (ER_VAR_DOES_NOT_EXIST)

  Message: Variable %s does not exist in persisted config file

  ER_VAR_DOES_NOT_EXIST was added in 8.0.1.

- Error: 3616 SQLSTATE: 22S02 (ER_LONGITUDE_OUT_OF_RANGE)

  Message: Longitude %f is out of range in function %s. It must be within (%f, %f].

  ER_LONGITUDE_OUT_OF_RANGE was added in 8.0.1.

- Error: 3617 SQLSTATE: 22S03 (ER_LATITUDE_OUT_OF_RANGE)

  Message: Latitude %f is out of range in function %s. It must be within [%f, %f].

  ER_LATITUDE_OUT_OF_RANGE was added in 8.0.1.

- Error: 3618 SQLSTATE: 22S00 (ER_NOT_IMPLEMENTED_FOR_GEOGRAPHIC_SRS)

  Message: %s(%s, %s) has not been implemented for geographic spatial reference systems.

  ER_NOT_IMPLEMENTED_FOR_GEOGRAPHIC_SRS was added in 8.0.1.

- Error: 3619 SQLSTATE: HY000 (ER_ILLEGAL_PRIVILEGE_LEVEL)

  Message: Illegal privilege level specified for %s

  ER_ILLEGAL_PRIVILEGE_LEVEL was added in 8.0.1.

- Error: 3620 SQLSTATE: HY000 (ER_NO_SYSTEM_VIEW_ACCESS)

  Message: Access to system view INFORMATION_SCHEMA.'%s' is rejected.

  ER_NO_SYSTEM_VIEW_ACCESS was added in 8.0.2.

- Error: 3621 SQLSTATE: XX999 (ER_PARSER_TRACE)

  Message: Parser saw: %s

  ER_PARSER_TRACE was added in 8.0.2.

- Error: 3622 SQLSTATE: HY000 (ER_BOOTSTRAP_CANT_THREAD)

  Message: Can't create thread to handle bootstrap (errno: %d)

  ER_BOOTSTRAP_CANT_THREAD was added in 8.0.2.

- Error: 3623 SQLSTATE: HY000 (ER_COMPONENT_FILTER_FLABBERGASTED)

  Message: The log-filter component "%s" got confused at "%s" ...

  ER_COMPONENT_FILTER_FLABBERGASTED was added in 8.0.2.

- Error: 3624 SQLSTATE: HY000 (ER_TRIGGER_INVALID_VALUE)

  Message: Trigger for table '%s'.'%s': invalid %s value (%s).

  ER_TRIGGER_INVALID_VALUE was added in 8.0.2.

- Error: 3625 SQLSTATE: HY000 (ER_OPT_WRONG_TREE)

  Message: Wrong tree: %s

  ER_OPT_WRONG_TREE was added in 8.0.2.

- Error: 3626 SQLSTATE: HY000 (ER_DD_FAILSAFE)

  Message: Error: Invalid %s

  ER_DD_FAILSAFE was added in 8.0.2.

- Error: 3627 SQLSTATE: HY000 (ER_DD_INIT_UPGRADE_FAILED)

  Message: Error in initializing dictionary, upgrade will do a cleanup and exit

  ER_DD_INIT_UPGRADE_FAILED was added in 8.0.2.

- Error: 3628 SQLSTATE: HY000 (ER_DD_NO_WRITES_NO_REPOPULATION)

  Message: Skip re-populating collations and character sets tables in %s%sread-only mode.

  ER_DD_NO_WRITES_NO_REPOPULATION was added in 8.0.2.

- Error: 3629 SQLSTATE: HY000 (ER_DD_VERSION_FOUND)

  Message: Found data dictionary with version %d

  ER_DD_VERSION_FOUND was added in 8.0.2.

- Error: 3630 SQLSTATE: HY000 (ER_DD_VERSION_INSTALLED)

  Message: Installed data dictionary with version %d

  ER_DD_VERSION_INSTALLED was added in 8.0.2.

- Error: 3631 SQLSTATE: HY000 (ER_DD_VERSION_UNSUPPORTED)

  Message: Data Dictionary version %d not supported

  ER_DD_VERSION_UNSUPPORTED was added in 8.0.2.

- Error: 3632 SQLSTATE: HY000 (ER_VIEW_UNKNOWN_CHARSET_OR_COLLATION)

  Message: View '%s'.'%s': unknown charset name and/or collation name (client: '%s'; connection: '%s').

  ER_VIEW_UNKNOWN_CHARSET_OR_COLLATION was added in 8.0.2.

- Error: 3633 SQLSTATE: HY000 (ER_DD_VIEW_CANT_ALLOC_CHARSET)

  Message: Error in allocating memory for character set name for view %s.%s.

  ER_DD_VIEW_CANT_ALLOC_CHARSET was added in 8.0.2.

- Error: 3634 SQLSTATE: HY000 (ER_DD_INIT_FAILED)

  Message: Data Dictionary initialization failed.

  ER_DD_INIT_FAILED was added in 8.0.2.

- Error: 3635 SQLSTATE: HY000 (ER_DD_UPDATING_PLUGIN_MD_FAILED)

  Message: Failed to update plugin metadata in dictionary tables.

  ER_DD_UPDATING_PLUGIN_MD_FAILED was added in 8.0.2.

- Error: 3636 SQLSTATE: HY000 (ER_DD_POPULATING_TABLES_FAILED)

  Message: Failed to Populate DD tables.

  ER_DD_POPULATING_TABLES_FAILED was added in 8.0.2.

- Error: 3637 SQLSTATE: HY000 (ER_DD_VIEW_CANT_CREATE)

  Message: Error in Creating View %s.%s

  ER_DD_VIEW_CANT_CREATE was added in 8.0.2.

- Error: 3638 SQLSTATE: HY000 (ER_DD_METADATA_NOT_FOUND)

  Message: Unable to start server. Cannot find the meta data for data dictionary table '%s'.

ER_DD_METADATA_NOT_FOUND was added in 8.0.2.

- Error: 3639 SQLSTATE: HY000 (ER_DD_CACHE_NOT_EMPTY_AT_SHUTDOWN)

  Message: Dictionary cache not empty at shutdown.

  ER_DD_CACHE_NOT_EMPTY_AT_SHUTDOWN was added in 8.0.2.

- Error: 3640 SQLSTATE: HY000 (ER_DD_OBJECT_REMAINS)

  Message: Dictionary objects used but not released.

  ER_DD_OBJECT_REMAINS was added in 8.0.2.

- Error: 3641 SQLSTATE: HY000 (ER_DD_OBJECT_REMAINS_IN_RELEASER)

  Message: Dictionary objects left in default releaser.

  ER_DD_OBJECT_REMAINS_IN_RELEASER was added in 8.0.2.

- Error: 3642 SQLSTATE: HY000 (ER_DD_OBJECT_RELEASER_REMAINS)

  Message: Dictionary object auto releaser not deleted

  ER_DD_OBJECT_RELEASER_REMAINS was added in 8.0.2.

- Error: 3643 SQLSTATE: HY000 (ER_DD_CANT_GET_OBJECT_KEY)

  Message: Error: Unable to create primary object key

  ER_DD_CANT_GET_OBJECT_KEY was added in 8.0.2.

- Error: 3644 SQLSTATE: HY000 (ER_DD_CANT_CREATE_OBJECT_KEY)

  Message: Error: Unable to create object key

  ER_DD_CANT_CREATE_OBJECT_KEY was added in 8.0.2.

- Error: 3645 SQLSTATE: HY000 (ER_LOG_SYSLOG_FACILITY_FAIL)

  Message: Failed to set syslog facility to "%s", setting to "%s" (%d) instead.

  ER_LOG_SYSLOG_FACILITY_FAIL was added in 8.0.2.

- Error: 3646 SQLSTATE: HY000 (ER_LOG_SYSLOG_CANNOT_OPEN)

  Message: Cannot open %s; check privileges, or start server with --log_syslog=0

  ER_LOG_SYSLOG_CANNOT_OPEN was added in 8.0.2.

- Error: 3647 SQLSTATE: HY000 (ER_LOG_FILE_CANNOT_OPEN)

  Message: Could not use %s for logging (error %d - %s). Turning logging off for the server process. To turn it on again: fix the cause, then%s restart the MySQL server.

  ER_LOG_FILE_CANNOT_OPEN was added in 8.0.2.

- Error: 3648 SQLSTATE: HY000 (ER_LOG_SLOW_CANNOT_OPEN)

  Message: either restart the query logging by using "SET GLOBAL SLOW_QUERY_LOG=ON" or

  ER_LOG_SLOW_CANNOT_OPEN was added in 8.0.2.

- Error: 3649 SQLSTATE: HY000 (ER_LOG_GENERAL_CANNOT_OPEN)

Message: either restart the query logging by using "SET GLOBAL GENERAL_LOG=ON" or

`ER_LOG_GENERAL_CANNOT_OPEN` was added in 8.0.2.

- Error: `3650` SQLSTATE: `HY000` (`ER_LOG_CANNOT_WRITE`)

  Message: Failed to write to %s: %s

  `ER_LOG_CANNOT_WRITE` was added in 8.0.2.

- Error: `3651` SQLSTATE: `HY000` (`ER_RPL_ZOMBIE_ENCOUNTERED`)

  Message: While initializing dump thread for slave with %s <%s>, found a zombie dump thread with the same %s. Master is killing the zombie dump thread(%u).

  `ER_RPL_ZOMBIE_ENCOUNTERED` was added in 8.0.2.

- Error: `3652` SQLSTATE: `HY000` (`ER_RPL_GTID_TABLE_CANNOT_OPEN`)

  Message: Gtid table is not ready to be used. Table '%s.%s' cannot be opened.

  `ER_RPL_GTID_TABLE_CANNOT_OPEN` was added in 8.0.2.

- Error: `3653` SQLSTATE: `HY000` (`ER_SYSTEM_SCHEMA_NOT_FOUND`)

  Message: System schema directory does not exist.

  `ER_SYSTEM_SCHEMA_NOT_FOUND` was added in 8.0.2.

- Error: `3654` SQLSTATE: `HY000` (`ER_CANT_CREATE_HANDLE_MGR_THREAD`)

  Message: Can't create handle_manager thread (errno= %d)

  `ER_CANT_CREATE_HANDLE_MGR_THREAD` was added in 8.0.2.

- Error: `3655` SQLSTATE: `HY000` (`ER_RPL_REPO_HAS_GAPS`)

  Message: It is not possible to change the type of the relay log's repository because there are workers' repositories with gaps. Please, fix the gaps first before doing such change.

  `ER_RPL_REPO_HAS_GAPS` was added in 8.0.2.

- Error: `3656` SQLSTATE: `HY000` (`ER_INVALID_VALUE_FOR_ENFORCE_GTID_CONSISTENCY`)

  Message: option 'enforce-gtid-consistency': value '%s' was not recognized. Setting enforce-gtid-consistency to OFF.

  `ER_INVALID_VALUE_FOR_ENFORCE_GTID_CONSISTENCY` was added in 8.0.2.

- Error: `3657` SQLSTATE: `HY000` (`ER_CHANGED_ENFORCE_GTID_CONSISTENCY`)

  Message: Changed ENFORCE_GTID_CONSISTENCY from %s to %s.

  `ER_CHANGED_ENFORCE_GTID_CONSISTENCY` was added in 8.0.2.

- Error: `3658` SQLSTATE: `HY000` (`ER_CHANGED_GTID_MODE`)

  Message: Changed GTID_MODE from %s to %s.

  `ER_CHANGED_GTID_MODE` was added in 8.0.2.

- Error: `3659` SQLSTATE: `HY000` (`ER_DISABLED_STORAGE_ENGINE_AS_DEFAULT`)

Message: %s is set to a disabled storage engine %s.

`ER_DISABLED_STORAGE_ENGINE_AS_DEFAULT` was added in 8.0.2.

- Error: `3660` SQLSTATE: `HY000` (`ER_DEBUG_SYNC_HIT`)

  Message: Debug sync points hit: %s

  `ER_DEBUG_SYNC_HIT` was added in 8.0.2.

- Error: `3661` SQLSTATE: `HY000` (`ER_DEBUG_SYNC_EXECUTED`)

  Message: Debug sync points executed: %s

  `ER_DEBUG_SYNC_EXECUTED` was added in 8.0.2.

- Error: `3662` SQLSTATE: `HY000` (`ER_DEBUG_SYNC_THREAD_MAX`)

  Message: Debug sync points max active per thread: %s

  `ER_DEBUG_SYNC_THREAD_MAX` was added in 8.0.2.

- Error: `3663` SQLSTATE: `HY000` (`ER_DEBUG_SYNC_OOM`)

  Message: Debug Sync Facility disabled due to lack of memory.

  `ER_DEBUG_SYNC_OOM` was added in 8.0.2.

- Error: `3664` SQLSTATE: `HY000` (`ER_CANT_INIT_TC_LOG`)

  Message: Can't init tc log

  `ER_CANT_INIT_TC_LOG` was added in 8.0.2.

- Error: `3665` SQLSTATE: `HY000` (`ER_EVENT_CANT_INIT_QUEUE`)

  Message: Event Scheduler: Can't initialize the execution queue

  `ER_EVENT_CANT_INIT_QUEUE` was added in 8.0.2.

- Error: `3666` SQLSTATE: `HY000` (`ER_EVENT_PURGING_QUEUE`)

  Message: Event Scheduler: Purging the queue. %u events

  `ER_EVENT_PURGING_QUEUE` was added in 8.0.2.

- Error: `3667` SQLSTATE: `HY000` (`ER_EVENT_LAST_EXECUTION`)

  Message: Event Scheduler: Last execution of %s.%s. %s

  `ER_EVENT_LAST_EXECUTION` was added in 8.0.2.

- Error: `3668` SQLSTATE: `HY000` (`ER_EVENT_MESSAGE_STACK`)

  Message: %*s

  `ER_EVENT_MESSAGE_STACK` was added in 8.0.2.

- Error: `3669` SQLSTATE: `HY000` (`ER_EVENT_EXECUTION_FAILED`)

  Message: Event Scheduler: [%s].[%s.%s] event execution failed.

  `ER_EVENT_EXECUTION_FAILED` was added in 8.0.2.

- Error: 3670 SQLSTATE: HY000 (ER_CANT_INIT_SCHEDULER_THREAD)

  Message: Event Scheduler: Cannot initialize the scheduler thread

  ER_CANT_INIT_SCHEDULER_THREAD was added in 8.0.2.

- Error: 3671 SQLSTATE: HY000 (ER_CANT_CREATE_SCHEDULER_THREAD)

  Message: Event scheduler: Failed to start scheduler, Can not create thread for event scheduler (errno=%d)

  ER_CANT_CREATE_SCHEDULER_THREAD was added in 8.0.2.

- Error: 3672 SQLSTATE: HY000 (ER_SCHEDULER_STOPPED)

  Message: Event Scheduler: Stopped

  ER_SCHEDULER_STOPPED was added in 8.0.2.

- Error: 3673 SQLSTATE: HY000 (ER_SCHEDULER_WAITING)

  Message: Event Scheduler: Waiting for the scheduler thread to reply

  ER_SCHEDULER_WAITING was added in 8.0.2.

- Error: 3674 SQLSTATE: HY000 (ER_SCHEDULER_STARTED)

  Message: Event Scheduler: scheduler thread started with id %u

  ER_SCHEDULER_STARTED was added in 8.0.2.

- Error: 3675 SQLSTATE: HY000 (ER_SCHEDULER_STOPPING_FAILED_TO_GET_EVENT)

  Message: Event Scheduler: Serious error during getting next event to execute. Stopping

  ER_SCHEDULER_STOPPING_FAILED_TO_GET_EVENT was added in 8.0.2.

- Error: 3676 SQLSTATE: HY000 (ER_SCHEDULER_STOPPING_FAILED_TO_CREATE_WORKER)

  Message: Event_scheduler::execute_top: Can not create event worker thread (errno=%d). Stopping event scheduler

  ER_SCHEDULER_STOPPING_FAILED_TO_CREATE_WORKER was added in 8.0.2.

- Error: 3677 SQLSTATE: HY000 (ER_SCHEDULER_KILLING)

  Message: Event Scheduler: Killing the scheduler thread, thread id %u

  ER_SCHEDULER_KILLING was added in 8.0.2.

- Error: 3678 SQLSTATE: HY000 (ER_UNABLE_TO_RESOLVE_IP)

  Message: IP address '%s' could not be resolved: %s

  ER_UNABLE_TO_RESOLVE_IP was added in 8.0.2.

- Error: 3679 SQLSTATE: HY000 (ER_UNABLE_TO_RESOLVE_HOSTNAME)

  Message: Host name '%s' could not be resolved: %s

  ER_UNABLE_TO_RESOLVE_HOSTNAME was added in 8.0.2.

- Error: 3680 SQLSTATE: HY000 (ER_HOSTNAME_RESEMBLES_IPV4)

Message: IP address '%s' has been resolved to the host name '%s', which resembles IPv4-address itself.

`ER_HOSTNAME_RESEMBLES_IPV4` was added in 8.0.2.

- Error: `3681` SQLSTATE: `HY000` (`ER_HOSTNAME_DOESNT_RESOLVE_TO`)

  Message: Hostname '%s' does not resolve to '%s'.

  `ER_HOSTNAME_DOESNT_RESOLVE_TO` was added in 8.0.2.

- Error: `3682` SQLSTATE: `HY000` (`ER_ADDRESSES_FOR_HOSTNAME_HEADER`)

  Message: Hostname '%s' has the following IP addresses:

  `ER_ADDRESSES_FOR_HOSTNAME_HEADER` was added in 8.0.2.

- Error: `3683` SQLSTATE: `HY000` (`ER_ADDRESSES_FOR_HOSTNAME_LIST_ITEM`)

  Message: - %s

  `ER_ADDRESSES_FOR_HOSTNAME_LIST_ITEM` was added in 8.0.2.

- Error: `3684` SQLSTATE: `HY000` (`ER_TRG_WITHOUT_DEFINER`)

  Message: Definer clause is missing in Trigger of Table %s. Rebuild Trigger to fix definer.

  `ER_TRG_WITHOUT_DEFINER` was added in 8.0.2.

- Error: `3685` SQLSTATE: `HY000` (`ER_TRG_NO_CLIENT_CHARSET`)

  Message: Client character set is missing for trigger of table %s. Using default character set.

  `ER_TRG_NO_CLIENT_CHARSET` was added in 8.0.2.

- Error: `3686` SQLSTATE: `HY000` (`ER_PARSING_VIEW`)

  Message: Error in parsing view %s.%s

  `ER_PARSING_VIEW` was added in 8.0.2.

- Error: `3687` SQLSTATE: `HY000` (`ER_COMPONENTS_INFRASTRUCTURE_BOOTSTRAP`)

  Message: Failed to bootstrap components infrastructure.

  `ER_COMPONENTS_INFRASTRUCTURE_BOOTSTRAP` was added in 8.0.2.

- Error: `3688` SQLSTATE: `HY000` (`ER_COMPONENTS_INFRASTRUCTURE_SHUTDOWN`)

  Message: Failed to shutdown components infrastructure.

  `ER_COMPONENTS_INFRASTRUCTURE_SHUTDOWN` was added in 8.0.2.

- Error: `3689` SQLSTATE: `HY000` (`ER_COMPONENTS_PERSIST_LOADER_BOOTSTRAP`)

  Message: Failed to bootstrap persistent components loader.

  `ER_COMPONENTS_PERSIST_LOADER_BOOTSTRAP` was added in 8.0.2.

- Error: `3690` SQLSTATE: `HY000` (`ER_DEPART_WITH_GRACE`)

  Message: Giving %d client threads a chance to die gracefully

  `ER_DEPART_WITH_GRACE` was added in 8.0.2.

- Error: 3691 SQLSTATE: HY000 (ER_CA_SELF_SIGNED)

  Message: CA certificate %s is self signed.

  ER_CA_SELF_SIGNED was added in 8.0.2.

- Error: 3692 SQLSTATE: HY000 (ER_SSL_LIBRARY_ERROR)

  Message: Failed to set up SSL because of the following SSL library error: %s

  ER_SSL_LIBRARY_ERROR was added in 8.0.2.

- Error: 3693 SQLSTATE: HY000 (ER_NO_THD_NO_UUID)

  Message: Failed to generate a server UUID because it is failed to allocate the THD.

  ER_NO_THD_NO_UUID was added in 8.0.2.

- Error: 3694 SQLSTATE: HY000 (ER_UUID_SALT)

  Message: Salting uuid generator variables, current_pid: %lu, server_start_time: %lu, bytes_sent: %llu,

  ER_UUID_SALT was added in 8.0.2.

- Error: 3695 SQLSTATE: HY000 (ER_UUID_IS)

  Message: Generated uuid: '%s', server_start_time: %lu, bytes_sent: %llu

  ER_UUID_IS was added in 8.0.2.

- Error: 3696 SQLSTATE: HY000 (ER_UUID_INVALID)

  Message: The server_uuid stored in auto.cnf file is not a valid UUID.

  ER_UUID_INVALID was added in 8.0.2.

- Error: 3697 SQLSTATE: HY000 (ER_UUID_SCRUB)

  Message: Garbage characters found at the end of the server_uuid value in auto.cnf file. It should be of length '%d' (UUID_LENGTH). Clear it and restart the server.

  ER_UUID_SCRUB was added in 8.0.2.

- Error: 3698 SQLSTATE: HY000 (ER_CREATING_NEW_UUID)

  Message: No existing UUID has been found, so we assume that this is the first time that this server has been started. Generating a new UUID: %s.

  ER_CREATING_NEW_UUID was added in 8.0.2.

- Error: 3699 SQLSTATE: HY000 (ER_CANT_CREATE_UUID)

  Message: Initialization of the server's UUID failed because it could not be read from the auto.cnf file. If this is a new server, the initialization failed because it was not possible to generate a new UUID.

  ER_CANT_CREATE_UUID was added in 8.0.2.

- Error: 3700 SQLSTATE: HY000 (ER_UNKNOWN_UNSUPPORTED_STORAGE_ENGINE)

  Message: Unknown/unsupported storage engine: %s

  ER_UNKNOWN_UNSUPPORTED_STORAGE_ENGINE was added in 8.0.2.

- Error: 3701 SQLSTATE: HY000 (ER_SECURE_AUTH_VALUE_UNSUPPORTED)

  Message: Unsupported value 0 for secure-auth

  ER_SECURE_AUTH_VALUE_UNSUPPORTED was added in 8.0.2.

- Error: 3702 SQLSTATE: HY000 (ER_INVALID_INSTRUMENT)

  Message: Invalid instrument name or value for performance_schema_instrument '%s'

  ER_INVALID_INSTRUMENT was added in 8.0.2.

- Error: 3703 SQLSTATE: HY000 (ER_INNODB_MANDATORY)

  Message: The use of InnoDB is mandatory since MySQL 5.7. The former options like '--innodb=0/1/OFF/ON' or '--skip-innodb' are ignored.

  ER_INNODB_MANDATORY was added in 8.0.2.

- Error: 3704 SQLSTATE: HY000 (ER_INNODB_CANNOT_BE_IGNORED)

  Message: ignore-builtin-innodb is ignored and will be removed in future releases.

  ER_INNODB_CANNOT_BE_IGNORED was added in 8.0.2.

- Error: 3705 SQLSTATE: HY000 (ER_OLD_PASSWORDS_NO_MIDDLE_GROUND)

  Message: Invalid old_passwords mode: 1. Valid values are 2 and 0

  ER_OLD_PASSWORDS_NO_MIDDLE_GROUND was added in 8.0.2.

- Error: 3706 SQLSTATE: HY000 (ER_VERBOSE_REQUIRES_HELP)

  Message: --verbose is for use with --help; did you mean --log-error-verbosity?

  ER_VERBOSE_REQUIRES_HELP was added in 8.0.2.

- Error: 3707 SQLSTATE: HY000 (ER_POINTLESS_WITHOUT_SLOWLOG)

  Message: options --log-slow-admin-statements, --log-queries-not-using-indexes and --log-slow-slave-statements have no effect if --slow-query-log is not set

  ER_POINTLESS_WITHOUT_SLOWLOG was added in 8.0.2.

- Error: 3708 SQLSTATE: HY000 (ER_WASTEFUL_NET_BUFFER_SIZE)

  Message: net_buffer_length (%lu) is set to be larger than max_allowed_packet (%lu). Please rectify.

  ER_WASTEFUL_NET_BUFFER_SIZE was added in 8.0.2.

- Error: 3709 SQLSTATE: HY000 (ER_DEPRECATED_TIMESTAMP_IMPLICIT_DEFAULTS)

  Message: TIMESTAMP with implicit DEFAULT value is deprecated. Please use --explicit_defaults_for_timestamp server option (see documentation for more details).

  ER_DEPRECATED_TIMESTAMP_IMPLICIT_DEFAULTS was added in 8.0.2.

- Error: 3710 SQLSTATE: HY000 (ER_FT_BOOL_SYNTAX_INVALID)

  Message: Invalid ft-boolean-syntax string: %s

  ER_FT_BOOL_SYNTAX_INVALID was added in 8.0.2.

- Error: 3711 SQLSTATE: HY000 (ER_CREDENTIALLESS_AUTO_USER_BAD)

Message: 'NO_AUTO_CREATE_USER' sql mode was not set.

ER_CREDENTIALLESS_AUTO_USER_BAD was added in 8.0.2.

- Error: 3712 SQLSTATE: HY000 (ER_CONNECTION_HANDLING_OOM)

  Message: Could not allocate memory for connection handling

  ER_CONNECTION_HANDLING_OOM was added in 8.0.2.

- Error: 3713 SQLSTATE: HY000 (ER_THREAD_HANDLING_OOM)

  Message: Could not allocate memory for thread handling

  ER_THREAD_HANDLING_OOM was added in 8.0.2.

- Error: 3714 SQLSTATE: HY000 (ER_CANT_CREATE_TEST_FILE)

  Message: Can't create test file %s

  ER_CANT_CREATE_TEST_FILE was added in 8.0.2.

- Error: 3715 SQLSTATE: HY000 (ER_CANT_CREATE_PID_FILE)

  Message: Can't start server: can't create PID file: %s

  ER_CANT_CREATE_PID_FILE was added in 8.0.2.

- Error: 3716 SQLSTATE: HY000 (ER_CANT_REMOVE_PID_FILE)

  Message: Unable to delete pid file: %s

  ER_CANT_REMOVE_PID_FILE was added in 8.0.2.

- Error: 3717 SQLSTATE: HY000 (ER_CANT_CREATE_SHUTDOWN_THREAD)

  Message: Can't create thread to handle shutdown requests (errno= %d)

  ER_CANT_CREATE_SHUTDOWN_THREAD was added in 8.0.2.

- Error: 3718 SQLSTATE: HY000 (ER_SEC_FILE_PRIV_CANT_ACCESS_DIR)

  Message: Failed to access directory for --secure-file-priv. Please make sure that directory exists and is accessible by MySQL Server. Supplied value : %s

  ER_SEC_FILE_PRIV_CANT_ACCESS_DIR was added in 8.0.2.

- Error: 3719 SQLSTATE: HY000 (ER_SEC_FILE_PRIV_IGNORED)

  Message: Ignoring --secure-file-priv value as server is running with --initialize(-insecure).

  ER_SEC_FILE_PRIV_IGNORED was added in 8.0.2.

- Error: 3720 SQLSTATE: HY000 (ER_SEC_FILE_PRIV_EMPTY)

  Message: Insecure configuration for --secure-file-priv: Current value does not restrict location of generated files. Consider setting it to a valid, non-empty path.

  ER_SEC_FILE_PRIV_EMPTY was added in 8.0.2.

- Error: 3721 SQLSTATE: HY000 (ER_SEC_FILE_PRIV_NULL)

Message: --secure-file-priv is set to NULL. Operations related to importing and exporting data are disabled

`ER_SEC_FILE_PRIV_NULL` was added in 8.0.2.

- Error: `3722` SQLSTATE: `HY000` (`ER_SEC_FILE_PRIV_DIRECTORY_INSECURE`)

  Message: Insecure configuration for --secure-file-priv: %s is accessible through --secure-file-priv. Consider choosing a different directory.

  `ER_SEC_FILE_PRIV_DIRECTORY_INSECURE` was added in 8.0.2.

- Error: `3723` SQLSTATE: `HY000` (`ER_SEC_FILE_PRIV_CANT_STAT`)

  Message: Failed to get stat for directory pointed out by --secure-file-priv

  `ER_SEC_FILE_PRIV_CANT_STAT` was added in 8.0.2.

- Error: `3724` SQLSTATE: `HY000` (`ER_SEC_FILE_PRIV_DIRECTORY_PERMISSIONS`)

  Message: Insecure configuration for --secure-file-priv: Location is accessible to all OS users. Consider choosing a different directory.

  `ER_SEC_FILE_PRIV_DIRECTORY_PERMISSIONS` was added in 8.0.2.

- Error: `3725` SQLSTATE: `HY000` (`ER_SEC_FILE_PRIV_ARGUMENT_TOO_LONG`)

  Message: Value for --secure-file-priv is longer than maximum limit of %d

  `ER_SEC_FILE_PRIV_ARGUMENT_TOO_LONG` was added in 8.0.2.

- Error: `3726` SQLSTATE: `HY000` (`ER_CANT_CREATE_NAMED_PIPES_THREAD`)

  Message: Can't create thread to handle named pipes (errno= %d)

  `ER_CANT_CREATE_NAMED_PIPES_THREAD` was added in 8.0.2.

- Error: `3727` SQLSTATE: `HY000` (`ER_CANT_CREATE_TCPIP_THREAD`)

  Message: Can't create thread to handle TCP/IP (errno= %d)

  `ER_CANT_CREATE_TCPIP_THREAD` was added in 8.0.2.

- Error: `3728` SQLSTATE: `HY000` (`ER_CANT_CREATE_SHM_THREAD`)

  Message: Can't create thread to handle shared memory (errno= %d)

  `ER_CANT_CREATE_SHM_THREAD` was added in 8.0.2.

- Error: `3729` SQLSTATE: `HY000` (`ER_CANT_CREATE_INTERRUPT_THREAD`)

  Message: Can't create interrupt-thread (error %d, errno: %d)

  `ER_CANT_CREATE_INTERRUPT_THREAD` was added in 8.0.2.

- Error: `3730` SQLSTATE: `HY000` (`ER_WRITABLE_CONFIG_REMOVED`)

  Message: World-writable config file '%s' has been removed.

  `ER_WRITABLE_CONFIG_REMOVED` was added in 8.0.2.

- Error: `3731` SQLSTATE: `HY000` (`ER_CORE_VALUES`)

Message: setrlimit could not change the size of core files to 'infinity'; We may not be able to generate a core file on signals

`ER_CORE_VALUES` was added in 8.0.2.

- Error: `3732` SQLSTATE: `HY000` (`ER_WRONG_DATETIME_SPEC`)

  Message: Wrong date/time format specifier: %s

  `ER_WRONG_DATETIME_SPEC` was added in 8.0.2.

- Error: `3733` SQLSTATE: `HY000` (`ER_RPL_BINLOG_FILTERS_OOM`)

  Message: Could not allocate replication and binlog filters: %s

  `ER_RPL_BINLOG_FILTERS_OOM` was added in 8.0.2.

- Error: `3734` SQLSTATE: `HY000` (`ER_KEYCACHE_OOM`)

  Message: Cannot allocate the keycache

  `ER_KEYCACHE_OOM` was added in 8.0.2.

- Error: `3735` SQLSTATE: `HY000` (`ER_CONFIRMING_THE_FUTURE`)

  Message: Current time has got past year 2038. Validating current time with %d iterations before initiating the normal server shutdown process.

  `ER_CONFIRMING_THE_FUTURE` was added in 8.0.2.

- Error: `3736` SQLSTATE: `HY000` (`ER_BACK_IN_TIME`)

  Message: Iteration %d: Obtained valid current time from system

  `ER_BACK_IN_TIME` was added in 8.0.2.

- Error: `3737` SQLSTATE: `HY000` (`ER_FUTURE_DATE`)

  Message: Iteration %d: Current time obtained from system is greater than 2038

  `ER_FUTURE_DATE` was added in 8.0.2.

- Error: `3738` SQLSTATE: `HY000` (`ER_UNSUPPORTED_DATE`)

  Message: This MySQL server doesn't support dates later then 2038

  `ER_UNSUPPORTED_DATE` was added in 8.0.2.

- Error: `3739` SQLSTATE: `HY000` (`ER_STARTING_AS`)

  Message: %s (mysqld %s) starting as process %lu ...

  `ER_STARTING_AS` was added in 8.0.2.

- Error: `3740` SQLSTATE: `HY000` (`ER_SHUTTING_DOWN_SLAVE_THREADS`)

  Message: Shutting down slave threads

  `ER_SHUTTING_DOWN_SLAVE_THREADS` was added in 8.0.2.

- Error: `3741` SQLSTATE: `HY000` (`ER_DISCONNECTING_REMAINING_CLIENTS`)

  Message: Forcefully disconnecting %d remaining clients

ER_DISCONNECTING_REMAINING_CLIENTS was added in 8.0.2.

- Error: 3742 SQLSTATE: HY000 (ER_ABORTING)

  Message: Aborting

  ER_ABORTING was added in 8.0.2.

- Error: 3743 SQLSTATE: HY000 (ER_BINLOG_END)

  Message: Binlog end

  ER_BINLOG_END was added in 8.0.2.

- Error: 3744 SQLSTATE: HY000 (ER_CALL_ME_LOCALHOST)

  Message: gethostname failed, using '%s' as hostname

  ER_CALL_ME_LOCALHOST was added in 8.0.2.

- Error: 3745 SQLSTATE: HY000 (ER_USER_REQUIRES_ROOT)

  Message: One can only use the --user switch if running as root

  ER_USER_REQUIRES_ROOT was added in 8.0.2.

- Error: 3746 SQLSTATE: HY000 (ER_REALLY_RUN_AS_ROOT)

  Message: Fatal error: Please read "Security" section of the manual to find out how to run mysqld as root!

  ER_REALLY_RUN_AS_ROOT was added in 8.0.2.

- Error: 3747 SQLSTATE: HY000 (ER_USER_WHAT_USER)

  Message: Fatal error: Can't change to run as user '%s' ; Please check that the user exists!

  ER_USER_WHAT_USER was added in 8.0.2.

- Error: 3748 SQLSTATE: HY000 (ER_TRANSPORTS_WHAT_TRANSPORTS)

  Message: Server is started with --require-secure-transport=ON but no secure transports (SSL or Shared Memory) are configured.

  ER_TRANSPORTS_WHAT_TRANSPORTS was added in 8.0.2.

- Error: 3749 SQLSTATE: HY000 (ER_FAIL_SETGID)

  Message: setgid: %s

  ER_FAIL_SETGID was added in 8.0.2.

- Error: 3750 SQLSTATE: HY000 (ER_FAIL_SETUID)

  Message: setuid: %s

  ER_FAIL_SETUID was added in 8.0.2.

- Error: 3751 SQLSTATE: HY000 (ER_FAIL_SETREGID)

  Message: setregid: %s

  ER_FAIL_SETREGID was added in 8.0.2.

- Error: 3752 SQLSTATE: HY000 (ER_FAIL_SETREUID)

  Message: setreuid: %s

  ER_FAIL_SETREUID was added in 8.0.2.

- Error: 3753 SQLSTATE: HY000 (ER_FAIL_CHROOT)

  Message: chroot: %s

  ER_FAIL_CHROOT was added in 8.0.2.

- Error: 3754 SQLSTATE: HY000 (ER_WIN_LISTEN_BUT_HOW)

  Message: TCP/IP, --shared-memory, or --named-pipe should be configured on NT OS

  ER_WIN_LISTEN_BUT_HOW was added in 8.0.2.

- Error: 3755 SQLSTATE: HY000 (ER_NOT_RIGHT_NOW)

  Message: CTRL-C ignored during startup

  ER_NOT_RIGHT_NOW was added in 8.0.2.

- Error: 3756 SQLSTATE: HY000 (ER_FIXING_CLIENT_CHARSET)

  Message: '%s' can not be used as client character set. '%s' will be used as default client character set.

  ER_FIXING_CLIENT_CHARSET was added in 8.0.2.

- Error: 3757 SQLSTATE: HY000 (ER_OOM)

  Message: Out of memory

  ER_OOM was added in 8.0.2.

- Error: 3758 SQLSTATE: HY000 (ER_FAILED_TO_LOCK_MEM)

  Message: Failed to lock memory. Errno: %d

  ER_FAILED_TO_LOCK_MEM was added in 8.0.2.

- Error: 3759 SQLSTATE: HY000 (ER_MYINIT_FAILED)

  Message: my_init() failed.

  ER_MYINIT_FAILED was added in 8.0.2.

- Error: 3760 SQLSTATE: HY000 (ER_BEG_INITFILE)

  Message: Execution of init_file \'%s\' started.

  ER_BEG_INITFILE was added in 8.0.2.

- Error: 3761 SQLSTATE: HY000 (ER_END_INITFILE)

  Message: Execution of init_file \'%s\' ended.

  ER_END_INITFILE was added in 8.0.2.

- Error: 3762 SQLSTATE: HY000 (ER_CHANGED_MAX_OPEN_FILES)

  Message: Changed limits: max_open_files: %lu (requested %lu)

ER_CHANGED_MAX_OPEN_FILES was added in 8.0.2.

- Error: 3763 SQLSTATE: HY000 (ER_CANT_INCREASE_MAX_OPEN_FILES)

  Message: Could not increase number of max_open_files to more than %lu (request: %lu)

  ER_CANT_INCREASE_MAX_OPEN_FILES was added in 8.0.2.

- Error: 3764 SQLSTATE: HY000 (ER_CHANGED_MAX_CONNECTIONS)

  Message: Changed limits: max_connections: %lu (requested %lu)

  ER_CHANGED_MAX_CONNECTIONS was added in 8.0.2.

- Error: 3765 SQLSTATE: HY000 (ER_CHANGED_TABLE_OPEN_CACHE)

  Message: Changed limits: table_open_cache: %lu (requested %lu)

  ER_CHANGED_TABLE_OPEN_CACHE was added in 8.0.2.

- Error: 3766 SQLSTATE: HY000 (ER_THE_USER_ABIDES)

  Message: Ignoring user change to '%s' because the user was set to '%s' earlier on the command line

  ER_THE_USER_ABIDES was added in 8.0.2.

- Error: 3767 SQLSTATE: HY000 (ER_RPL_CANT_ADD_DO_TABLE)

  Message: Could not add do table rule '%s'!

  ER_RPL_CANT_ADD_DO_TABLE was added in 8.0.2.

- Error: 3768 SQLSTATE: HY000 (ER_RPL_CANT_ADD_IGNORE_TABLE)

  Message: Could not add ignore table rule '%s'!

  ER_RPL_CANT_ADD_IGNORE_TABLE was added in 8.0.2.

- Error: 3769 SQLSTATE: HY000 (ER_TRACK_VARIABLES_BOGUS)

  Message: The variable session_track_system_variables either has duplicate values or invalid values.

  ER_TRACK_VARIABLES_BOGUS was added in 8.0.2.

- Error: 3770 SQLSTATE: HY000 (ER_EXCESS_ARGUMENTS)

  Message: Too many arguments (first extra is '%s').

  ER_EXCESS_ARGUMENTS was added in 8.0.2.

- Error: 3771 SQLSTATE: HY000 (ER_VERBOSE_HINT)

  Message: Use --verbose --help to get a list of available options!

  ER_VERBOSE_HINT was added in 8.0.2.

- Error: 3772 SQLSTATE: HY000 (ER_CANT_READ_ERRMSGS)

  Message: Unable to read errmsg.sys file

  ER_CANT_READ_ERRMSGS was added in 8.0.2.

- Error: 3773 SQLSTATE: HY000 (ER_CANT_INIT_DBS)

Message: Can't init databases

`ER_CANT_INIT_DBS` was added in 8.0.2.

- Error: `3774` SQLSTATE: `HY000` (`ER_LOG_OUTPUT_CONTRADICTORY`)

  Message: There were other values specified to log-output besides NONE. Disabling slow and general logs anyway.

  `ER_LOG_OUTPUT_CONTRADICTORY` was added in 8.0.2.

- Error: `3775` SQLSTATE: `HY000` (`ER_NO_CSV_NO_LOG_TABLES`)

  Message: CSV engine is not present, falling back to the log files

  `ER_NO_CSV_NO_LOG_TABLES` was added in 8.0.2.

- Error: `3776` SQLSTATE: `HY000` (`ER_RPL_REWRITEDB_MISSING_ARROW`)

  Message: Bad syntax in replicate-rewrite-db - missing '->'!

  `ER_RPL_REWRITEDB_MISSING_ARROW` was added in 8.0.2.

- Error: `3777` SQLSTATE: `HY000` (`ER_RPL_REWRITEDB_EMPTY_FROM`)

  Message: Bad syntax in replicate-rewrite-db - empty FROM db!

  `ER_RPL_REWRITEDB_EMPTY_FROM` was added in 8.0.2.

- Error: `3778` SQLSTATE: `HY000` (`ER_RPL_REWRITEDB_EMPTY_TO`)

  Message: Bad syntax in replicate-rewrite-db - empty TO db!

  `ER_RPL_REWRITEDB_EMPTY_TO` was added in 8.0.2.

- Error: `3779` SQLSTATE: `HY000` (`ER_LOG_FILES_GIVEN_LOG_OUTPUT_IS_TABLE`)

  Message: Although a path was specified for the %s, log tables are used. To enable logging to files use the --log-output=file option.

  `ER_LOG_FILES_GIVEN_LOG_OUTPUT_IS_TABLE` was added in 8.0.2.

- Error: `3780` SQLSTATE: `HY000` (`ER_LOG_FILE_INVALID`)

  Message: Invalid value for %s: %s

  `ER_LOG_FILE_INVALID` was added in 8.0.2.

- Error: `3781` SQLSTATE: `HY000` (`ER_LOWER_CASE_TABLE_NAMES_CS_DD_ON_CI_FS_UNSUPPORTED`)

  Message: The server option 'lower_case_table_names' is configured to use case sensitive table names but the data directory is on a case-insensitive file system which is an unsupported combination. Please consider either using a case sensitive file system for your data directory or switching to a case-insensitive table name mode.

  `ER_LOWER_CASE_TABLE_NAMES_CS_DD_ON_CI_FS_UNSUPPORTED` was added in 8.0.2.

- Error: `3782` SQLSTATE: `HY000` (`ER_LOWER_CASE_TABLE_NAMES_USING_2`)

  Message: Setting lower_case_table_names=2 because file system for %s is case insensitive

  `ER_LOWER_CASE_TABLE_NAMES_USING_2` was added in 8.0.2.

- Error: 3783 SQLSTATE: HY000 (ER_LOWER_CASE_TABLE_NAMES_USING_0)

  Message: lower_case_table_names was set to 2, even though your the file system '%s' is case sensitive. Now setting lower_case_table_names to 0 to avoid future problems.

  ER_LOWER_CASE_TABLE_NAMES_USING_0 was added in 8.0.2.

- Error: 3784 SQLSTATE: HY000 (ER_NEED_LOG_BIN)

  Message: You need to use --log-bin to make %s work.

  ER_NEED_LOG_BIN was added in 8.0.2.

- Error: 3785 SQLSTATE: HY000 (ER_NEED_FILE_INSTEAD_OF_DIR)

  Message: Path '%s' is a directory name, please specify a file name for %s option

  ER_NEED_FILE_INSTEAD_OF_DIR was added in 8.0.2.

- Error: 3786 SQLSTATE: HY000 (ER_LOG_BIN_BETTER_WITH_NAME)

  Message: No argument was provided to --log-bin, and --log-bin-index was not used; so replication may break when this MySQL server acts as a master and has his hostname changed!! Please use '--log-bin=%s' to avoid this problem.

  ER_LOG_BIN_BETTER_WITH_NAME was added in 8.0.2.

- Error: 3787 SQLSTATE: HY000 (ER_BINLOG_NEEDS_SERVERID)

  Message: You have enabled the binary log, but you haven't provided the mandatory server-id. Please refer to the proper server start-up parameters documentation

  ER_BINLOG_NEEDS_SERVERID was added in 8.0.2.

- Error: 3788 SQLSTATE: HY000 (ER_RPL_CANT_MAKE_PATHS)

  Message: Unable to create replication path names: out of memory or path names too long (path name exceeds %d or file name exceeds %d).

  ER_RPL_CANT_MAKE_PATHS was added in 8.0.2.

- Error: 3789 SQLSTATE: HY000 (ER_CANT_INITIALIZE_GTID)

  Message: Failed to initialize GTID structures.

  ER_CANT_INITIALIZE_GTID was added in 8.0.2.

- Error: 3790 SQLSTATE: HY000 (ER_CANT_INITIALIZE_EARLY_PLUGINS)

  Message: Failed to initialize early plugins.

  ER_CANT_INITIALIZE_EARLY_PLUGINS was added in 8.0.2.

- Error: 3791 SQLSTATE: HY000 (ER_CANT_INITIALIZE_BUILTIN_PLUGINS)

  Message: Failed to initialize builtin plugins.

  ER_CANT_INITIALIZE_BUILTIN_PLUGINS was added in 8.0.2.

- Error: 3792 SQLSTATE: HY000 (ER_CANT_INITIALIZE_DYNAMIC_PLUGINS)

  Message: Failed to initialize dynamic plugins.

  ER_CANT_INITIALIZE_DYNAMIC_PLUGINS was added in 8.0.2.

- Error: 3793 SQLSTATE: HY000 (ER_PERFSCHEMA_INIT_FAILED)

  Message: Performance schema disabled (reason: init failed).

  ER_PERFSCHEMA_INIT_FAILED was added in 8.0.2.

- Error: 3794 SQLSTATE: HY000 (ER_STACKSIZE_UNEXPECTED)

  Message: Asked for %lu thread stack, but got %ld

  ER_STACKSIZE_UNEXPECTED was added in 8.0.2.

- Error: 3795 SQLSTATE: HY000 (ER_CANT_SET_DATADIR)

  Message: failed to set datadir to %s

  ER_CANT_SET_DATADIR was added in 8.0.2.

- Error: 3796 SQLSTATE: HY000 (ER_CANT_STAT_DATADIR)

  Message: Can't read data directory's stats (%d): %s. Assuming that it's not owned by the same user/group

  ER_CANT_STAT_DATADIR was added in 8.0.2.

- Error: 3797 SQLSTATE: HY000 (ER_CANT_CHOWN_DATADIR)

  Message: Can't change data directory owner to %s

  ER_CANT_CHOWN_DATADIR was added in 8.0.2.

- Error: 3798 SQLSTATE: HY000 (ER_CANT_SET_UP_PERSISTED_VALUES)

  Message: Setting persistent options failed.

  ER_CANT_SET_UP_PERSISTED_VALUES was added in 8.0.2.

- Error: 3799 SQLSTATE: HY000 (ER_CANT_SAVE_GTIDS)

  Message: Failed to save the set of Global Transaction Identifiers of the last binary log into the mysql.gtid_executed table while the server was shutting down. The next server restart will make another attempt to save Global Transaction Identifiers into the table.

  ER_CANT_SAVE_GTIDS was added in 8.0.2.

- Error: 3800 SQLSTATE: HY000 (ER_AUTH_CANT_SET_DEFAULT_PLUGIN)

  Message: Can't start server: Invalid value for --default-authentication-plugin

  ER_AUTH_CANT_SET_DEFAULT_PLUGIN was added in 8.0.2.

- Error: 3801 SQLSTATE: HY000 (ER_CANT_JOIN_SHUTDOWN_THREAD)

  Message: Could not join %sthread. error:%d

  ER_CANT_JOIN_SHUTDOWN_THREAD was added in 8.0.2.

- Error: 3802 SQLSTATE: HY000 (ER_CANT_HASH_DO_AND_IGNORE_RULES)

  Message: An error occurred while building do_table and ignore_table rules to hashes for global replication filter.

  ER_CANT_HASH_DO_AND_IGNORE_RULES was added in 8.0.2.

- Error: 3803 SQLSTATE: HY000 (ER_CANT_OPEN_CA)

  Message: Error opening CA certificate file

  ER_CANT_OPEN_CA was added in 8.0.2.

- Error: 3804 SQLSTATE: HY000 (ER_CANT_ACCESS_CAPATH)

  Message: Error accessing directory pointed by --ssl-capath

  ER_CANT_ACCESS_CAPATH was added in 8.0.2.

- Error: 3805 SQLSTATE: HY000 (ER_SSL_TRYING_DATADIR_DEFAULTS)

  Message: Found %s, %s and %s in data directory. Trying to enable SSL support using them.

  ER_SSL_TRYING_DATADIR_DEFAULTS was added in 8.0.2.

- Error: 3806 SQLSTATE: HY000 (ER_AUTO_OPTIONS_FAILED)

  Message: Failed to create %s(file: '%s', errno %d)

  ER_AUTO_OPTIONS_FAILED was added in 8.0.2.

- Error: 3807 SQLSTATE: HY000 (ER_CANT_INIT_TIMER)

  Message: Failed to initialize timer component (errno %d).

  ER_CANT_INIT_TIMER was added in 8.0.2.

- Error: 3808 SQLSTATE: HY000 (ER_SERVERID_TOO_LARGE)

  Message: server-id configured is too large to represent with server-id-bits configured.

  ER_SERVERID_TOO_LARGE was added in 8.0.2.

- Error: 3809 SQLSTATE: HY000 (ER_DEFAULT_SE_UNAVAILABLE)

  Message: Default%s storage engine (%s) is not available

  ER_DEFAULT_SE_UNAVAILABLE was added in 8.0.2.

- Error: 3810 SQLSTATE: HY000 (ER_CANT_OPEN_ERROR_LOG)

  Message: Could not open %s file for error logging: %s.

  ER_CANT_OPEN_ERROR_LOG was added in 8.0.2.

- Error: 3811 SQLSTATE: HY000 (ER_INVALID_ERROR_LOG_NAME)

  Message: Invalid log file name after expanding symlinks: '%s'

  ER_INVALID_ERROR_LOG_NAME was added in 8.0.2.

- Error: 3812 SQLSTATE: HY000 (ER_RPL_INFINITY_DENIED)

  Message: using --replicate-same-server-id in conjunction with --log-slave-updates is impossible, it would lead to infinite loops in this server.

  ER_RPL_INFINITY_DENIED was added in 8.0.2.

- Error: 3813 SQLSTATE: HY000 (ER_RPL_INFINITY_IGNORED)

Message: using --replicate-same-server-id in conjunction with --log-slave-updates would lead to infinite loops in this server. However this will be ignored as the --log-bin option is not defined.

`ER_RPL_INFINITY_IGNORED` was added in 8.0.2.

- Error: `3814` SQLSTATE: `HY000` (`ER_NDB_TABLES_NOT_READY`)

  Message: NDB : Tables not available after %lu seconds. Consider increasing --ndb-wait-setup value

  `ER_NDB_TABLES_NOT_READY` was added in 8.0.2.

- Error: `3815` SQLSTATE: `HY000` (`ER_TABLE_CHECK_INTACT`)

  Message: %s

  `ER_TABLE_CHECK_INTACT` was added in 8.0.2.

- Error: `3816` SQLSTATE: `HY000` (`ER_DD_TABLESPACE_NOT_FOUND`)

  Message: Unable to start server. The data dictionary tablespace '%s' does not exist.

  `ER_DD_TABLESPACE_NOT_FOUND` was added in 8.0.2.

- Error: `3817` SQLSTATE: `HY000` (`ER_DD_TRG_CONNECTION_COLLATION_MISSING`)

  Message: Connection collation is missing for trigger of table %s. Using default connection collation.

  `ER_DD_TRG_CONNECTION_COLLATION_MISSING` was added in 8.0.2.

- Error: `3818` SQLSTATE: `HY000` (`ER_DD_TRG_DB_COLLATION_MISSING`)

  Message: Database collation is missing for trigger of table %s. Using Default character set.

  `ER_DD_TRG_DB_COLLATION_MISSING` was added in 8.0.2.

- Error: `3819` SQLSTATE: `HY000` (`ER_DD_TRG_DEFINER_OOM`)

  Message: Error in Memory allocation for Definer %s for Trigger.

  `ER_DD_TRG_DEFINER_OOM` was added in 8.0.2.

- Error: `3820` SQLSTATE: `HY000` (`ER_DD_TRG_FILE_UNREADABLE`)

  Message: Error in reading %s.TRG file.

  `ER_DD_TRG_FILE_UNREADABLE` was added in 8.0.2.

- Error: `3821` SQLSTATE: `HY000` (`ER_TRG_CANT_PARSE`)

  Message: Error in parsing Triggers from %s.TRG file.

  `ER_TRG_CANT_PARSE` was added in 8.0.2.

- Error: `3822` SQLSTATE: `HY000` (`ER_DD_TRG_CANT_ADD`)

  Message: Error in creating DD entry for Trigger %s.%s

  `ER_DD_TRG_CANT_ADD` was added in 8.0.2.

- Error: `3823` SQLSTATE: `HY000` (`ER_DD_CANT_RESOLVE_VIEW`)

  Message: Resolving dependency for the view '%s.%s' failed. View is no more valid to use

  `ER_DD_CANT_RESOLVE_VIEW` was added in 8.0.2.

- Error: 3824 SQLSTATE: HY000 (ER_DD_VIEW_WITHOUT_DEFINER)

  Message: %s.%s has no definer (as per an old view format). Current user is used as definer. Please recreate the view.

  ER_DD_VIEW_WITHOUT_DEFINER was added in 8.0.2.

- Error: 3825 SQLSTATE: HY000 (ER_PLUGIN_INIT_FAILED)

  Message: Plugin '%s' init function returned error.

  ER_PLUGIN_INIT_FAILED was added in 8.0.2.

- Error: 3826 SQLSTATE: HY000 (ER_RPL_TRX_DELEGATES_INIT_FAILED)

  Message: Initialization of transaction delegates failed. Please report a bug.

  ER_RPL_TRX_DELEGATES_INIT_FAILED was added in 8.0.2.

- Error: 3827 SQLSTATE: HY000 (ER_RPL_BINLOG_STORAGE_DELEGATES_INIT_FAILED)

  Message: Initialization binlog storage delegates failed. Please report a bug.

  ER_RPL_BINLOG_STORAGE_DELEGATES_INIT_FAILED was added in 8.0.2.

- Error: 3828 SQLSTATE: HY000 (ER_RPL_BINLOG_TRANSMIT_DELEGATES_INIT_FAILED)

  Message: Initialization of binlog transmit delegates failed. Please report a bug.

  ER_RPL_BINLOG_TRANSMIT_DELEGATES_INIT_FAILED was added in 8.0.2.

- Error: 3829 SQLSTATE: HY000 (ER_RPL_BINLOG_RELAY_DELEGATES_INIT_FAILED)

  Message: Initialization binlog relay IO delegates failed. Please report a bug.

  ER_RPL_BINLOG_RELAY_DELEGATES_INIT_FAILED was added in 8.0.2.

- Error: 3830 SQLSTATE: HY000 (ER_RPL_PLUGIN_FUNCTION_FAILED)

  Message: Run function '...' in plugin '%s' failed

  ER_RPL_PLUGIN_FUNCTION_FAILED was added in 8.0.2.

- Error: 3831 SQLSTATE: HY000 (ER_SQL_HA_READ_FAILED)

  Message: mysql_ha_read: Got error %d when reading table '%s'

  ER_SQL_HA_READ_FAILED was added in 8.0.2.

- Error: 3832 SQLSTATE: HY000 (ER_SR_BOGUS_VALUE)

  Message: Stored routine '%s'.'%s': invalid value in column %s.

  ER_SR_BOGUS_VALUE was added in 8.0.2.

- Error: 3833 SQLSTATE: HY000 (ER_SR_INVALID_CONTEXT)

  Message: Invalid creation context '%s.%s'.

  ER_SR_INVALID_CONTEXT was added in 8.0.2.

- Error: 3834 SQLSTATE: HY000 (ER_READING_TABLE_FAILED)

  Message: Got error %d when reading table '%s'

ER_READING_TABLE_FAILED was added in 8.0.2.

- Error: 3835 SQLSTATE: HY000 (ER_DES_FILE_WRONG_KEY)

  Message: load_des_file: Found wrong key_number: %c

  ER_DES_FILE_WRONG_KEY was added in 8.0.2.

- Error: 3836 SQLSTATE: HY000 (ER_CANT_SET_PERSISTED)

  Message: Failed to set persisted options.

  ER_CANT_SET_PERSISTED was added in 8.0.2.

- Error: 3837 SQLSTATE: HY000 (ER_JSON_PARSE_ERROR)

  Message: JSON parsing error

  ER_JSON_PARSE_ERROR was added in 8.0.2.

- Error: 3838 SQLSTATE: HY000 (ER_CONFIG_OPTION_WITHOUT_GROUP)

  Message: Found option without preceding group in config file

  ER_CONFIG_OPTION_WITHOUT_GROUP was added in 8.0.2.

- Error: 3839 SQLSTATE: HY000 (ER_VALGRIND_DO_QUICK_LEAK_CHECK)

  Message: VALGRIND_DO_QUICK_LEAK_CHECK

  ER_VALGRIND_DO_QUICK_LEAK_CHECK was added in 8.0.2.

- Error: 3840 SQLSTATE: HY000 (ER_VALGRIND_COUNT_LEAKS)

  Message: VALGRIND_COUNT_LEAKS reports %lu leaked bytes for query '%.*s'

  ER_VALGRIND_COUNT_LEAKS was added in 8.0.2.

- Error: 3841 SQLSTATE: HY000 (ER_LOAD_DATA_INFILE_FAILED_IN_UNEXPECTED_WAY)

  Message: LOAD DATA INFILE in the slave SQL Thread can only read from --slave-load-tmpdir. Please, report a bug.

  ER_LOAD_DATA_INFILE_FAILED_IN_UNEXPECTED_WAY was added in 8.0.2.

- Error: 3842 SQLSTATE: HY000 (ER_UNKNOWN_ERROR_NUMBER)

  Message: Got unknown error: %d

  ER_UNKNOWN_ERROR_NUMBER was added in 8.0.2.

- Error: 3843 SQLSTATE: HY000 (ER_UDF_CANT_ALLOC_FOR_STRUCTURES)

  Message: Can't allocate memory for udf structures

  ER_UDF_CANT_ALLOC_FOR_STRUCTURES was added in 8.0.2.

- Error: 3844 SQLSTATE: HY000 (ER_UDF_CANT_ALLOC_FOR_FUNCTION)

  Message: Can't alloc memory for udf function: '%s'

  ER_UDF_CANT_ALLOC_FOR_FUNCTION was added in 8.0.2.

- Error: 3845 SQLSTATE: HY000 (ER_UDF_INVALID_ROW_IN_FUNCTION_TABLE)

Message: Invalid row in mysql.func table for function '%s'

ER_UDF_INVALID_ROW_IN_FUNCTION_TABLE was added in 8.0.2.

- Error: 3846 SQLSTATE: HY000 (ER_UDF_CANT_OPEN_FUNCTION_TABLE)

  Message: Can't open the mysql.func table. Please run mysql_upgrade to create it.

  ER_UDF_CANT_OPEN_FUNCTION_TABLE was added in 8.0.2.

- Error: 3847 SQLSTATE: HY000 (ER_XA_RECOVER_FOUND_TRX_IN_SE)

  Message: Found %d prepared transaction(s) in %s

  ER_XA_RECOVER_FOUND_TRX_IN_SE was added in 8.0.2.

- Error: 3848 SQLSTATE: HY000 (ER_XA_RECOVER_FOUND_XA_TRX)

  Message: Found %d prepared XA transactions

  ER_XA_RECOVER_FOUND_XA_TRX was added in 8.0.2.

- Error: 3849 SQLSTATE: HY000 (ER_XA_IGNORING_XID)

  Message: ignore xid %s

  ER_XA_IGNORING_XID was added in 8.0.2.

- Error: 3850 SQLSTATE: HY000 (ER_XA_COMMITTING_XID)

  Message: commit xid %s

  ER_XA_COMMITTING_XID was added in 8.0.2.

- Error: 3851 SQLSTATE: HY000 (ER_XA_ROLLING_BACK_XID)

  Message: rollback xid %s

  ER_XA_ROLLING_BACK_XID was added in 8.0.2.

- Error: 3852 SQLSTATE: HY000 (ER_XA_STARTING_RECOVERY)

  Message: Starting crash recovery...

  ER_XA_STARTING_RECOVERY was added in 8.0.2.

- Error: 3853 SQLSTATE: HY000 (ER_XA_NO_MULTI_2PC_HEURISTIC_RECOVER)

  Message: --tc-heuristic-recover rollback strategy is not safe on systems with more than one 2-phase-commit-capable storage engine. Aborting crash recovery.

  ER_XA_NO_MULTI_2PC_HEURISTIC_RECOVER was added in 8.0.2.

- Error: 3854 SQLSTATE: HY000 (ER_XA_RECOVER_EXPLANATION)

  Message: Found %d prepared transactions! It means that mysqld was not shut down properly last time and critical recovery information (last binlog or %s file) was manually deleted after a crash. You have to start mysqld with --tc-heuristic-recover switch to commit or rollback pending transactions.

  ER_XA_RECOVER_EXPLANATION was added in 8.0.2.

- Error: 3855 SQLSTATE: HY000 (ER_XA_RECOVERY_DONE)

Message: Crash recovery finished.

`ER_XA_RECOVERY_DONE` was added in 8.0.2.

- Error: `3856` SQLSTATE: `HY000` (`ER_TRX_GTID_COLLECT_REJECT`)

  Message: Failed to collect GTID to send in the response packet!

  `ER_TRX_GTID_COLLECT_REJECT` was added in 8.0.2.

- Error: `3857` SQLSTATE: `HY000` (`ER_SQL_AUTHOR_DEFAULT_ROLES_FAIL`)

  Message: MYSQL.DEFAULT_ROLES couldn't be updated for authorizationidentifier %s

  `ER_SQL_AUTHOR_DEFAULT_ROLES_FAIL` was added in 8.0.2.

- Error: `3858` SQLSTATE: `HY000` (`ER_SQL_USER_TABLE_CREATE_WARNING`)

  Message: Following users were specified in CREATE USER IF NOT EXISTS but they already exist. Corresponding entry in binary log used default authentication plugin '%s' to rewrite authentication information(if any) for them: %s

  `ER_SQL_USER_TABLE_CREATE_WARNING` was added in 8.0.2.

- Error: `3859` SQLSTATE: `HY000` (`ER_SQL_USER_TABLE_ALTER_WARNING`)

  Message: Following users were specified in ALTER USER IF EXISTS but they do not exist. Corresponding entry in binary log used default authentication plugin '%s' to rewrite authentication information(if any) for them: %s

  `ER_SQL_USER_TABLE_ALTER_WARNING` was added in 8.0.2.

- Error: `3860` SQLSTATE: `HY000` (`ER_ROW_IN_WRONG_PARTITION_PLEASE_REPAIR`)

  Message: Table '%s' corrupted: row in wrong partition: %s -- Please REPAIR the table!

  `ER_ROW_IN_WRONG_PARTITION_PLEASE_REPAIR` was added in 8.0.2.

- Error: `3861` SQLSTATE: `HY000` (`ER_MYISAM_CRASHED_ERROR_IN_THREAD`)

  Message: Got an error from thread_id=%u, %s:%d

  `ER_MYISAM_CRASHED_ERROR_IN_THREAD` was added in 8.0.2.

- Error: `3862` SQLSTATE: `HY000` (`ER_MYISAM_CRASHED_ERROR_IN`)

  Message: Got an error from unknown thread, %s:%d

  `ER_MYISAM_CRASHED_ERROR_IN` was added in 8.0.2.

- Error: `3863` SQLSTATE: `HY000` (`ER_TOO_MANY_STORAGE_ENGINES`)

  Message: Too many storage engines!

  `ER_TOO_MANY_STORAGE_ENGINES` was added in 8.0.2.

- Error: `3864` SQLSTATE: `HY000` (`ER_SE_TYPECODE_CONFLICT`)

  Message: Storage engine '%s' has conflicting typecode. Assigning value %d.

  `ER_SE_TYPECODE_CONFLICT` was added in 8.0.2.

- Error: `3865` SQLSTATE: `HY000` (`ER_TRX_WRITE_SET_OOM`)

Message: Out of memory on transaction write set extraction

`ER_TRX_WRITE_SET_OOM` was added in 8.0.2.

- Error: `3866` SQLSTATE: `HY000` (`ER_HANDLERTON_OOM`)

  Message: Unable to allocate memory for plugin '%s' handlerton.

  `ER_HANDLERTON_OOM` was added in 8.0.2.

- Error: `3867` SQLSTATE: `HY000` (`ER_CONN_SHM_LISTENER`)

  Message: Shared memory setting up listener

  `ER_CONN_SHM_LISTENER` was added in 8.0.2.

- Error: `3868` SQLSTATE: `HY000` (`ER_CONN_SHM_CANT_CREATE_SERVICE`)

  Message: Can't create shared memory service: %s. : %s

  `ER_CONN_SHM_CANT_CREATE_SERVICE` was added in 8.0.2.

- Error: `3869` SQLSTATE: `HY000` (`ER_CONN_SHM_CANT_CREATE_CONNECTION`)

  Message: Can't create shared memory connection: %s. : %s

  `ER_CONN_SHM_CANT_CREATE_CONNECTION` was added in 8.0.2.

- Error: `3870` SQLSTATE: `HY000` (`ER_CONN_PIP_CANT_CREATE_EVENT`)

  Message: Can't create event, last error=%u

  `ER_CONN_PIP_CANT_CREATE_EVENT` was added in 8.0.2.

- Error: `3871` SQLSTATE: `HY000` (`ER_CONN_PIP_CANT_CREATE_PIPE`)

  Message: Can't create new named pipe!: %s

  `ER_CONN_PIP_CANT_CREATE_PIPE` was added in 8.0.2.

- Error: `3872` SQLSTATE: `HY000` (`ER_CONN_PER_THREAD_NO_THREAD`)

  Message: Can't create thread to handle new connection(errno= %d)

  `ER_CONN_PER_THREAD_NO_THREAD` was added in 8.0.2.

- Error: `3873` SQLSTATE: `HY000` (`ER_CONN_TCP_NO_SOCKET`)

  Message: Failed to create a socket for %s '%s': errno: %d.

  `ER_CONN_TCP_NO_SOCKET` was added in 8.0.2.

- Error: `3874` SQLSTATE: `HY000` (`ER_CONN_TCP_CREATED`)

  Message: Server socket created on IP: '%s'.

  `ER_CONN_TCP_CREATED` was added in 8.0.2.

- Error: `3875` SQLSTATE: `HY000` (`ER_CONN_TCP_ADDRESS`)

  Message: Server hostname (bind-address): '%s'; port: %d

  `ER_CONN_TCP_ADDRESS` was added in 8.0.2.

- Error: 3876 SQLSTATE: HY000 (ER_CONN_TCP_IPV6_AVAILABLE)

  Message: IPv6 is available.

  ER_CONN_TCP_IPV6_AVAILABLE was added in 8.0.2.

- Error: 3877 SQLSTATE: HY000 (ER_CONN_TCP_IPV6_UNAVAILABLE)

  Message: IPv6 is not available.

  ER_CONN_TCP_IPV6_UNAVAILABLE was added in 8.0.2.

- Error: 3878 SQLSTATE: HY000 (ER_CONN_TCP_ERROR_WITH_STRERROR)

  Message: Can't create IP socket: %s

  ER_CONN_TCP_ERROR_WITH_STRERROR was added in 8.0.2.

- Error: 3879 SQLSTATE: HY000 (ER_CONN_TCP_CANT_RESOLVE_HOSTNAME)

  Message: Can't start server: cannot resolve hostname!

  ER_CONN_TCP_CANT_RESOLVE_HOSTNAME was added in 8.0.2.

- Error: 3880 SQLSTATE: HY000 (ER_CONN_TCP_IS_THERE_ANOTHER_USING_PORT)

  Message: Do you already have another mysqld server running on port: %d ?

  ER_CONN_TCP_IS_THERE_ANOTHER_USING_PORT was added in 8.0.2.

- Error: 3881 SQLSTATE: HY000 (ER_CONN_UNIX_IS_THERE_ANOTHER_USING_SOCKET)

  Message: Do you already have another mysqld server running on socket: %s ?

  ER_CONN_UNIX_IS_THERE_ANOTHER_USING_SOCKET was added in 8.0.2.

- Error: 3882 SQLSTATE: HY000 (ER_CONN_UNIX_PID_CLAIMED_SOCKET_FILE)

  Message: Another process with pid %d is using unix socket file.

  ER_CONN_UNIX_PID_CLAIMED_SOCKET_FILE was added in 8.0.2.

- Error: 3883 SQLSTATE: HY000 (ER_CONN_TCP_CANT_RESET_V6ONLY)

  Message: Failed to reset IPV6_V6ONLY flag (error: %d). The server will listen to IPv6 addresses only.

  ER_CONN_TCP_CANT_RESET_V6ONLY was added in 8.0.2.

- Error: 3884 SQLSTATE: HY000 (ER_CONN_TCP_BIND_RETRY)

  Message: Retrying bind on TCP/IP port %u

  ER_CONN_TCP_BIND_RETRY was added in 8.0.2.

- Error: 3885 SQLSTATE: HY000 (ER_CONN_TPC_BIND_FAIL)

  Message: Can't start server: Bind on TCP/IP port: %s

  ER_CONN_TPC_BIND_FAIL was added in 8.0.2.

- Error: 3886 SQLSTATE: HY000 (ER_CONN_TCP_IP_NOT_LOGGED)

  Message: Fails to print out IP-address.

ER_CONN_TCP_IP_NOT_LOGGED was added in 8.0.2.

- Error: 3887 SQLSTATE: HY000 (ER_CONN_TCP_RESOLVE_INFO)

  Message: - '%s' resolves to '%s';

  ER_CONN_TCP_RESOLVE_INFO was added in 8.0.2.

- Error: 3888 SQLSTATE: HY000 (ER_CONN_TCP_START_FAIL)

  Message: Can't start server: listen() on TCP/IP port: %s

  ER_CONN_TCP_START_FAIL was added in 8.0.2.

- Error: 3889 SQLSTATE: HY000 (ER_CONN_TCP_LISTEN_FAIL)

  Message: listen() on TCP/IP failed with error %d

  ER_CONN_TCP_LISTEN_FAIL was added in 8.0.2.

- Error: 3890 SQLSTATE: HY000 (ER_CONN_UNIX_PATH_TOO_LONG)

  Message: The socket file path is too long (> %u): %s

  ER_CONN_UNIX_PATH_TOO_LONG was added in 8.0.2.

- Error: 3891 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_FAIL)

  Message: Unable to setup unix socket lock file.

  ER_CONN_UNIX_LOCK_FILE_FAIL was added in 8.0.2.

- Error: 3892 SQLSTATE: HY000 (ER_CONN_UNIX_NO_FD)

  Message: Can't start server: UNIX Socket : %s

  ER_CONN_UNIX_NO_FD was added in 8.0.2.

- Error: 3893 SQLSTATE: HY000 (ER_CONN_UNIX_NO_BIND_NO_START)

  Message: Can't start server : Bind on unix socket: %s

  ER_CONN_UNIX_NO_BIND_NO_START was added in 8.0.2.

- Error: 3894 SQLSTATE: HY000 (ER_CONN_UNIX_LISTEN_FAILED)

  Message: listen() on Unix socket failed with error %d

  ER_CONN_UNIX_LISTEN_FAILED was added in 8.0.2.

- Error: 3895 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_GIVING_UP)

  Message: Unable to create unix socket lock file %s after retries.

  ER_CONN_UNIX_LOCK_FILE_GIVING_UP was added in 8.0.2.

- Error: 3896 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_CREATE)

  Message: Could not create unix socket lock file %s.

  ER_CONN_UNIX_LOCK_FILE_CANT_CREATE was added in 8.0.2.

- Error: 3897 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_OPEN)

Message: Could not open unix socket lock file %s.

ER_CONN_UNIX_LOCK_FILE_CANT_OPEN was added in 8.0.2.

- Error: 3898 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_READ)

  Message: Could not read unix socket lock file %s.

  ER_CONN_UNIX_LOCK_FILE_CANT_READ was added in 8.0.2.

- Error: 3899 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_EMPTY)

  Message: Unix socket lock file is empty %s.

  ER_CONN_UNIX_LOCK_FILE_EMPTY was added in 8.0.2.

- Error: 3900 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_PIDLESS)

  Message: Invalid pid in unix socket lock file %s.

  ER_CONN_UNIX_LOCK_FILE_PIDLESS was added in 8.0.2.

- Error: 3901 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_WRITE)

  Message: Could not write unix socket lock file %s errno %d.

  ER_CONN_UNIX_LOCK_FILE_CANT_WRITE was added in 8.0.2.

- Error: 3902 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_DELETE)

  Message: Could not remove unix socket lock file %s errno %d.

  ER_CONN_UNIX_LOCK_FILE_CANT_DELETE was added in 8.0.2.

- Error: 3903 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_SYNC)

  Message: Could not sync unix socket lock file %s errno %d.

  ER_CONN_UNIX_LOCK_FILE_CANT_SYNC was added in 8.0.2.

- Error: 3904 SQLSTATE: HY000 (ER_CONN_UNIX_LOCK_FILE_CANT_CLOSE)

  Message: Could not close unix socket lock file %s errno %d.

  ER_CONN_UNIX_LOCK_FILE_CANT_CLOSE was added in 8.0.2.

- Error: 3905 SQLSTATE: HY000 (ER_CONN_SOCKET_SELECT_FAILED)

  Message: mysqld: Got error %d from select

  ER_CONN_SOCKET_SELECT_FAILED was added in 8.0.2.

- Error: 3906 SQLSTATE: HY000 (ER_CONN_SOCKET_ACCEPT_FAILED)

  Message: Error in accept: %s

  ER_CONN_SOCKET_ACCEPT_FAILED was added in 8.0.2.

- Error: 3907 SQLSTATE: HY000 (ER_AUTH_RSA_CANT_FIND)

  Message: RSA %s key file not found: %s. Some authentication plugins will not work.

  ER_AUTH_RSA_CANT_FIND was added in 8.0.2.

- Error: 3908 SQLSTATE: HY000 (ER_AUTH_RSA_CANT_PARSE)

  Message: Failure to parse RSA %s key (file exists): %s: %s

  ER_AUTH_RSA_CANT_PARSE was added in 8.0.2.

- Error: 3909 SQLSTATE: HY000 (ER_AUTH_RSA_CANT_READ)

  Message: Failure to read key file: %s

  ER_AUTH_RSA_CANT_READ was added in 8.0.2.

- Error: 3910 SQLSTATE: HY000 (ER_AUTH_RSA_FILES_NOT_FOUND)

  Message: RSA key files not found. Some authentication plugins will not work.

  ER_AUTH_RSA_FILES_NOT_FOUND was added in 8.0.2.

- Error: 3911 SQLSTATE: HY000 (ER_CONN_ATTR_TRUNCATED)

  Message: Connection attributes of length %lu were truncated (%d bytes lost) for connection %llu, user %s@%s (as %s), auth: %s

  ER_CONN_ATTR_TRUNCATED was added in 8.0.2.

- Error: 3912 SQLSTATE: HY000 (ER_X509_CIPHERS_MISMATCH)

  Message: X509 ciphers mismatch: should be '%s' but is '%s'

  ER_X509_CIPHERS_MISMATCH was added in 8.0.2.

- Error: 3913 SQLSTATE: HY000 (ER_X509_ISSUER_MISMATCH)

  Message: X509 issuer mismatch: should be '%s' but is '%s'

  ER_X509_ISSUER_MISMATCH was added in 8.0.2.

- Error: 3914 SQLSTATE: HY000 (ER_X509_SUBJECT_MISMATCH)

  Message: X509 subject mismatch: should be '%s' but is '%s'

  ER_X509_SUBJECT_MISMATCH was added in 8.0.2.

- Error: 3915 SQLSTATE: HY000 (ER_AUTH_CANT_ACTIVATE_ROLE)

  Message: Failed to activate default role %s for %s

  ER_AUTH_CANT_ACTIVATE_ROLE was added in 8.0.2.

- Error: 3916 SQLSTATE: HY000 (ER_X509_NEEDS_RSA_PRIVKEY)

  Message: Could not generate RSA private key required for X509 certificate.

  ER_X509_NEEDS_RSA_PRIVKEY was added in 8.0.2.

- Error: 3917 SQLSTATE: HY000 (ER_X509_CANT_WRITE_KEY)

  Message: Could not write key file: %s

  ER_X509_CANT_WRITE_KEY was added in 8.0.2.

- Error: 3918 SQLSTATE: HY000 (ER_X509_CANT_CHMOD_KEY)

  Message: Could not set file permission for %s

ER_X509_CANT_CHMOD_KEY was added in 8.0.2.

- Error: 3919 SQLSTATE: HY000 (ER_X509_CANT_READ_CA_KEY)

  Message: Could not read CA key file: %s

  ER_X509_CANT_READ_CA_KEY was added in 8.0.2.

- Error: 3920 SQLSTATE: HY000 (ER_X509_CANT_READ_CA_CERT)

  Message: Could not read CA certificate file: %s

  ER_X509_CANT_READ_CA_CERT was added in 8.0.2.

- Error: 3921 SQLSTATE: HY000 (ER_X509_CANT_CREATE_CERT)

  Message: Could not generate X509 certificate.

  ER_X509_CANT_CREATE_CERT was added in 8.0.2.

- Error: 3922 SQLSTATE: HY000 (ER_X509_CANT_WRITE_CERT)

  Message: Could not write certificate file: %s

  ER_X509_CANT_WRITE_CERT was added in 8.0.2.

- Error: 3923 SQLSTATE: HY000 (ER_AUTH_CANT_CREATE_RSA_PAIR)

  Message: Could not generate RSA Private/Public key pair

  ER_AUTH_CANT_CREATE_RSA_PAIR was added in 8.0.2.

- Error: 3924 SQLSTATE: HY000 (ER_AUTH_CANT_WRITE_PRIVKEY)

  Message: Could not write private key file: %s

  ER_AUTH_CANT_WRITE_PRIVKEY was added in 8.0.2.

- Error: 3925 SQLSTATE: HY000 (ER_AUTH_CANT_WRITE_PUBKEY)

  Message: Could not write public key file: %s

  ER_AUTH_CANT_WRITE_PUBKEY was added in 8.0.2.

- Error: 3926 SQLSTATE: HY000 (ER_AUTH_SSL_CONF_PREVENTS_CERT_GENERATION)

  Message: Skipping generation of SSL certificates as options related to SSL are specified.

  ER_AUTH_SSL_CONF_PREVENTS_CERT_GENERATION was added in 8.0.2.

- Error: 3927 SQLSTATE: HY000 (ER_AUTH_USING_EXISTING_CERTS)

  Message: Skipping generation of SSL certificates as certificate files are present in data directory.

  ER_AUTH_USING_EXISTING_CERTS was added in 8.0.2.

- Error: 3928 SQLSTATE: HY000 (ER_AUTH_CERTS_SAVED_TO_DATADIR)

  Message: Auto generated SSL certificates are placed in data directory.

  ER_AUTH_CERTS_SAVED_TO_DATADIR was added in 8.0.2.

- Error: 3929 SQLSTATE: HY000 (ER_AUTH_CERT_GENERATION_DISABLED)

Message: Skipping generation of SSL certificates as --auto_generate_certs is set to OFF.

ER_AUTH_CERT_GENERATION_DISABLED was added in 8.0.2.

- Error: 3930 SQLSTATE: HY000 (ER_AUTH_RSA_CONF_PREVENTS_KEY_GENERATION)

  Message: Skipping generation of RSA key pair as options related to RSA keys are specified.

  ER_AUTH_RSA_CONF_PREVENTS_KEY_GENERATION was added in 8.0.2.

- Error: 3931 SQLSTATE: HY000 (ER_AUTH_KEY_GENERATION_SKIPPED_PAIR_PRESENT)

  Message: Skipping generation of RSA key pair as key files are present in data directory.

  ER_AUTH_KEY_GENERATION_SKIPPED_PAIR_PRESENT was added in 8.0.2.

- Error: 3932 SQLSTATE: HY000 (ER_AUTH_KEYS_SAVED_TO_DATADIR)

  Message: Auto generated RSA key files are placed in data directory.

  ER_AUTH_KEYS_SAVED_TO_DATADIR was added in 8.0.2.

- Error: 3933 SQLSTATE: HY000 (ER_AUTH_KEY_GENERATION_DISABLED)

  Message: Skipping generation of RSA key pair as --sha256_password_auto_generate_rsa_keys is set to OFF.

  ER_AUTH_KEY_GENERATION_DISABLED was added in 8.0.2.

- Error: 3934 SQLSTATE: HY000 (ER_AUTHCACHE_PROXIES_PRIV_SKIPPED_NEEDS_RESOLVE)

  Message: 'proxies_priv' entry '%s@%s %s@%s' ignored in --skip-name-resolve mode.

  ER_AUTHCACHE_PROXIES_PRIV_SKIPPED_NEEDS_RESOLVE was added in 8.0.2.

- Error: 3935 SQLSTATE: HY000 (ER_AUTHCACHE_PLUGIN_MISSING)

  Message: The plugin '%.*s' used to authenticate user '%s'@'%.*s' is not loaded. Nobody can currently login using this account.

  ER_AUTHCACHE_PLUGIN_MISSING was added in 8.0.2.

- Error: 3936 SQLSTATE: HY000 (ER_AUTHCACHE_PLUGIN_CONFIG)

  Message: The plugin '%s' is used to authenticate user '%s'@'%.*s', %s configured. Nobody can currently login using this account.

  ER_AUTHCACHE_PLUGIN_CONFIG was added in 8.0.2.

- Error: 3937 SQLSTATE: HY000 (ER_AUTHCACHE_ROLE_TABLES_DODGY)

  Message: Could not load mysql.role_edges and mysql.default_roles tables. ACL DDLs will not work unless mysql_upgrade is executed.

  ER_AUTHCACHE_ROLE_TABLES_DODGY was added in 8.0.2.

- Error: 3938 SQLSTATE: HY000 (ER_AUTHCACHE_USER_SKIPPED_NEEDS_RESOLVE)

  Message: 'user' entry '%s@%s' ignored in --skip-name-resolve mode.

  ER_AUTHCACHE_USER_SKIPPED_NEEDS_RESOLVE was added in 8.0.2.

- Error: 3939 SQLSTATE: HY000 (ER_AUTHCACHE_USER_TABLE_DODGY)

Message: Fatal error: mysql.user table is damaged. Please run mysql_upgrade.

ER_AUTHCACHE_USER_TABLE_DODGY was added in 8.0.2.

- Error: 3940 SQLSTATE: HY000 (ER_AUTHCACHE_USER_IGNORED_DEPRECATED_PASSWORD)

  Message: User entry '%s'@'%s' has a deprecated pre-4.1 password. The user will be ignored and no one can login with this user anymore.

  ER_AUTHCACHE_USER_IGNORED_DEPRECATED_PASSWORD was added in 8.0.2.

- Error: 3941 SQLSTATE: HY000 (ER_AUTHCACHE_USER_IGNORED_NEEDS_PLUGIN)

  Message: User entry '%s'@'%s' has an empty plugin value. The user will be ignored and no one can login with this user anymore.

  ER_AUTHCACHE_USER_IGNORED_NEEDS_PLUGIN was added in 8.0.2.

- Error: 3942 SQLSTATE: HY000 (ER_AUTHCACHE_USER_IGNORED_INVALID_PASSWORD)

  Message: Found invalid password for user: '%s@%s'; Ignoring user

  ER_AUTHCACHE_USER_IGNORED_INVALID_PASSWORD was added in 8.0.2.

- Error: 3943 SQLSTATE: HY000 (ER_AUTHCACHE_EXPIRED_PASSWORD_UNSUPPORTED)

  Message: 'user' entry '%s@%s' has the password ignore flag raised, but its authentication plugin doesn't support password expiration. The user id will be ignored.

  ER_AUTHCACHE_EXPIRED_PASSWORD_UNSUPPORTED was added in 8.0.2.

- Error: 3944 SQLSTATE: HY000 (ER_NO_SUPER_WITHOUT_USER_PLUGIN)

  Message: Some of the user accounts with SUPER privileges were disabled because of empty mysql.user.plugin value. If you are upgrading from MySQL 5.6 to MySQL 5.7 it means we were not able to substitute for empty plugin column. Probably because of pre 4.1 password hash. If your account is disabled you will need to: 1. Stop the server and restart it with --skip-grant-tables. 2. Run mysql_upgrade. 3. Restart the server with the parameters you normally use. For complete instructions on how to upgrade MySQL to a new version please see the 'Upgrading MySQL' section from the MySQL manual

  ER_NO_SUPER_WITHOUT_USER_PLUGIN was added in 8.0.2.

- Error: 3945 SQLSTATE: HY000 (ER_AUTHCACHE_DB_IGNORED_EMPTY_NAME)

  Message: Found an entry in the 'db' table with empty database name; Skipped

  ER_AUTHCACHE_DB_IGNORED_EMPTY_NAME was added in 8.0.2.

- Error: 3946 SQLSTATE: HY000 (ER_AUTHCACHE_DB_SKIPPED_NEEDS_RESOLVE)

  Message: 'db' entry '%s %s@%s' ignored in --skip-name-resolve mode.

  ER_AUTHCACHE_DB_SKIPPED_NEEDS_RESOLVE was added in 8.0.2.

- Error: 3947 SQLSTATE: HY000 (ER_AUTHCACHE_DB_ENTRY_LOWERCASED_REVOKE_WILL_FAIL)

  Message: 'db' entry '%s %s@%s' had database in mixed case that has been forced to lowercase because lower_case_table_names is set. It will not be possible to remove this privilege using REVOKE.

ER_AUTHCACHE_DB_ENTRY_LOWERCASED_REVOKE_WILL_FAIL was added in 8.0.2.

- Error: 3948 SQLSTATE: HY000 (ER_AUTHCACHE_TABLE_PROXIES_PRIV_MISSING)

  Message: Missing system table mysql.proxies_priv; please run mysql_upgrade to create it

  ER_AUTHCACHE_TABLE_PROXIES_PRIV_MISSING was added in 8.0.2.

- Error: 3949 SQLSTATE: HY000 (ER_AUTHCACHE_CANT_OPEN_AND_LOCK_PRIVILEGE_TABLES)

  Message: Fatal error: Can't open and lock privilege tables: %s

  ER_AUTHCACHE_CANT_OPEN_AND_LOCK_PRIVILEGE_TABLES was added in 8.0.2.

- Error: 3950 SQLSTATE: HY000 (ER_AUTHCACHE_CANT_INIT_GRANT_SUBSYSTEM)

  Message: Fatal: can't initialize grant subsystem - '%s'

  ER_AUTHCACHE_CANT_INIT_GRANT_SUBSYSTEM was added in 8.0.2.

- Error: 3951 SQLSTATE: HY000 (ER_AUTHCACHE_PROCS_PRIV_SKIPPED_NEEDS_RESOLVE)

  Message: 'procs_priv' entry '%s %s@%s' ignored in --skip-name-resolve mode.

  ER_AUTHCACHE_PROCS_PRIV_SKIPPED_NEEDS_RESOLVE was added in 8.0.2.

- Error: 3952 SQLSTATE: HY000 (ER_AUTHCACHE_PROCS_PRIV_ENTRY_IGNORED_BAD_ROUTINE_TYPE)

  Message: 'procs_priv' entry '%s' ignored, bad routine type

  ER_AUTHCACHE_PROCS_PRIV_ENTRY_IGNORED_BAD_ROUTINE_TYPE was added in 8.0.2.

- Error: 3953 SQLSTATE: HY000 (ER_AUTHCACHE_TABLES_PRIV_SKIPPED_NEEDS_RESOLVE)

  Message: 'tables_priv' entry '%s %s@%s' ignored in --skip-name-resolve mode.

  ER_AUTHCACHE_TABLES_PRIV_SKIPPED_NEEDS_RESOLVE was added in 8.0.2.

- Error: 3954 SQLSTATE: HY000 (ER_USER_NOT_IN_EXTRA_USERS_BINLOG_POSSIBLY_INCOMPLETE)

  Message: Failed to add %s in extra_users. Binary log entry may miss some of the users.

  ER_USER_NOT_IN_EXTRA_USERS_BINLOG_POSSIBLY_INCOMPLETE was added in 8.0.2.

- Error: 3955 SQLSTATE: HY000 (ER_DD_SCHEMA_NOT_FOUND)

  Message: Unable to start server. The data dictionary schema '%s' does not exist.

  ER_DD_SCHEMA_NOT_FOUND was added in 8.0.2.

- Error: 3956 SQLSTATE: HY000 (ER_DD_TABLE_NOT_FOUND)

  Message: Unable to start server. The data dictionary table '%s' does not exist.

  ER_DD_TABLE_NOT_FOUND was added in 8.0.2.

- Error: 3957 SQLSTATE: HY000 (ER_DD_SE_INIT_FAILED)

  Message: Failed to initialize DD Storage Engine

  ER_DD_SE_INIT_FAILED was added in 8.0.2.

- Error: 3958 SQLSTATE: HY000 (ER_DD_ABORTING_PARTIAL_UPGRADE)

  Message: Found partially upgraded DD. Aborting upgrade and deleting all DD tables. Start the upgrade process again.

  ER_DD_ABORTING_PARTIAL_UPGRADE was added in 8.0.2.

- Error: 3959 SQLSTATE: HY000 (ER_DD_FRM_EXISTS_FOR_TABLE)

  Message: Found .frm file with same name as one of the Dictionary Tables.

  ER_DD_FRM_EXISTS_FOR_TABLE was added in 8.0.2.

- Error: 3960 SQLSTATE: HY000 (ER_DD_CREATED_FOR_UPGRADE)

  Message: Created Data Dictionary for upgrade

  ER_DD_CREATED_FOR_UPGRADE was added in 8.0.2.

- Error: 3961 SQLSTATE: HY000 (ER_ERRMSG_CANT_FIND_FILE)

  Message: Can't find error-message file '%s'. Check error-message file location and 'lc-messages-dir' configuration directive.

  ER_ERRMSG_CANT_FIND_FILE was added in 8.0.2.

- Error: 3962 SQLSTATE: HY000 (ER_ERRMSG_LOADING_55_STYLE)

  Message: Using pre 5.5 semantics to load error messages from %s. If this is not intended, refer to the documentation for valid usage of --lc-messages-dir and --language parameters.

  ER_ERRMSG_LOADING_55_STYLE was added in 8.0.2.

- Error: 3963 SQLSTATE: HY000 (ER_ERRMSG_MISSING_IN_FILE)

  Message: Error message file '%s' had only %d error messages, but it should contain at least %d error messages. Check that the above file is the right version for this program!

  ER_ERRMSG_MISSING_IN_FILE was added in 8.0.2.

- Error: 3964 SQLSTATE: HY000 (ER_ERRMSG_OOM)

  Message: Not enough memory for messagefile '%s'

  ER_ERRMSG_OOM was added in 8.0.2.

- Error: 3965 SQLSTATE: HY000 (ER_ERRMSG_CANT_READ)

  Message: Can't read from messagefile '%s'

  ER_ERRMSG_CANT_READ was added in 8.0.2.

- Error: 3966 SQLSTATE: HY000 (ER_TABLE_INCOMPATIBLE_DECIMAL_FIELD)

  Message: Found incompatible DECIMAL field '%s' in %s; Please do "ALTER TABLE `%s` FORCE" to fix it!

  ER_TABLE_INCOMPATIBLE_DECIMAL_FIELD was added in 8.0.2.

- Error: 3967 SQLSTATE: HY000 (ER_TABLE_INCOMPATIBLE_YEAR_FIELD)

  Message: Found incompatible YEAR(x) field '%s' in %s; Please do "ALTER TABLE `%s` FORCE" to fix it!

ER_TABLE_INCOMPATIBLE_YEAR_FIELD was added in 8.0.2.

- Error: 3968 SQLSTATE: HY000 (ER_INVALID_CHARSET_AND_DEFAULT_IS_MB)

  Message: '%s' had no or invalid character set, and default character set is multi-byte, so character column sizes may have changed

  ER_INVALID_CHARSET_AND_DEFAULT_IS_MB was added in 8.0.2.

- Error: 3969 SQLSTATE: HY000 (ER_TABLE_WRONG_KEY_DEFINITION)

  Message: Found wrong key definition in %s; Please do "ALTER TABLE `%s` FORCE " to fix it!

  ER_TABLE_WRONG_KEY_DEFINITION was added in 8.0.2.

- Error: 3970 SQLSTATE: HY000 (ER_CANT_OPEN_FRM_FILE)

  Message: Unable to open file %s

  ER_CANT_OPEN_FRM_FILE was added in 8.0.2.

- Error: 3971 SQLSTATE: HY000 (ER_CANT_READ_FRM_FILE)

  Message: Error in reading file %s

  ER_CANT_READ_FRM_FILE was added in 8.0.2.

- Error: 3972 SQLSTATE: HY000 (ER_TABLE_CREATED_WITH_DIFFERENT_VERSION)

  Message: Table '%s' was created with a different version of MySQL and cannot be read

  ER_TABLE_CREATED_WITH_DIFFERENT_VERSION was added in 8.0.2.

- Error: 3973 SQLSTATE: HY000 (ER_VIEW_UNPARSABLE)

  Message: Unable to read view %s

  ER_VIEW_UNPARSABLE was added in 8.0.2.

- Error: 3974 SQLSTATE: HY000 (ER_FILE_TYPE_UNKNOWN)

  Message: File %s has unknown type in its header.

  ER_FILE_TYPE_UNKNOWN was added in 8.0.2.

- Error: 3975 SQLSTATE: HY000 (ER_INVALID_INFO_IN_FRM)

  Message: Incorrect information in file %s

  ER_INVALID_INFO_IN_FRM was added in 8.0.2.

- Error: 3976 SQLSTATE: HY000 (ER_CANT_OPEN_AND_LOCK_PRIVILEGE_TABLES)

  Message: Can't open and lock privilege tables: %s

  ER_CANT_OPEN_AND_LOCK_PRIVILEGE_TABLES was added in 8.0.2.

- Error: 3977 SQLSTATE: HY000 (ER_AUDIT_PLUGIN_DOES_NOT_SUPPORT_AUDIT_AUTH_EVENTS)

  Message: Plugin '%s' cannot subscribe to MYSQL_AUDIT_AUTHORIZATION events. Currently not supported.

ER_AUDIT_PLUGIN_DOES_NOT_SUPPORT_AUDIT_AUTH_EVENTS was added in 8.0.2.

- Error: 3978 SQLSTATE: HY000 (ER_AUDIT_PLUGIN_HAS_INVALID_DATA)

  Message: Plugin '%s' has invalid data.

  ER_AUDIT_PLUGIN_HAS_INVALID_DATA was added in 8.0.2.

- Error: 3979 SQLSTATE: HY000 (ER_TZ_OOM_INITIALIZING_TIME_ZONES)

  Message: Fatal error: OOM while initializing time zones

  ER_TZ_OOM_INITIALIZING_TIME_ZONES was added in 8.0.2.

- Error: 3980 SQLSTATE: HY000 (ER_TZ_CANT_OPEN_AND_LOCK_TIME_ZONE_TABLE)

  Message: Can't open and lock time zone table: %s trying to live without them

  ER_TZ_CANT_OPEN_AND_LOCK_TIME_ZONE_TABLE was added in 8.0.2.

- Error: 3981 SQLSTATE: HY000 (ER_TZ_OOM_LOADING_LEAP_SECOND_TABLE)

  Message: Fatal error: Out of memory while loading mysql.time_zone_leap_second table

  ER_TZ_OOM_LOADING_LEAP_SECOND_TABLE was added in 8.0.2.

- Error: 3982 SQLSTATE: HY000 (ER_TZ_TOO_MANY_LEAPS_IN_LEAP_SECOND_TABLE)

  Message: Fatal error: While loading mysql.time_zone_leap_second table: too much leaps

  ER_TZ_TOO_MANY_LEAPS_IN_LEAP_SECOND_TABLE was added in 8.0.2.

- Error: 3983 SQLSTATE: HY000 (ER_TZ_ERROR_LOADING_LEAP_SECOND_TABLE)

  Message: Fatal error: Error while loading mysql.time_zone_leap_second table

  ER_TZ_ERROR_LOADING_LEAP_SECOND_TABLE was added in 8.0.2.

- Error: 3984 SQLSTATE: HY000 (ER_TZ_UNKNOWN_OR_ILLEGAL_DEFAULT_TIME_ZONE)

  Message: Fatal error: Illegal or unknown default time zone '%s'

  ER_TZ_UNKNOWN_OR_ILLEGAL_DEFAULT_TIME_ZONE was added in 8.0.2.

- Error: 3985 SQLSTATE: HY000 (ER_TZ_CANT_FIND_DESCRIPTION_FOR_TIME_ZONE)

  Message: Can't find description of time zone '%.*s'

  ER_TZ_CANT_FIND_DESCRIPTION_FOR_TIME_ZONE was added in 8.0.2.

- Error: 3986 SQLSTATE: HY000 (ER_TZ_CANT_FIND_DESCRIPTION_FOR_TIME_ZONE_ID)

  Message: Can't find description of time zone '%u'

  ER_TZ_CANT_FIND_DESCRIPTION_FOR_TIME_ZONE_ID was added in 8.0.2.

- Error: 3987 SQLSTATE: HY000 (ER_TZ_TRANSITION_TYPE_TABLE_TYPE_TOO_LARGE)

  Message: Error while loading time zone description from mysql.time_zone_transition_type table: too big transition type id

  ER_TZ_TRANSITION_TYPE_TABLE_TYPE_TOO_LARGE was added in 8.0.2.

- Error: 3988 SQLSTATE: HY000 (ER_TZ_TRANSITION_TYPE_TABLE_ABBREVIATIONS_EXCEED_SPACE)

  Message: Error while loading time zone description from mysql.time_zone_transition_type table: not enough room for abbreviations

  ER_TZ_TRANSITION_TYPE_TABLE_ABBREVIATIONS_EXCEED_SPACE was added in 8.0.2.

- Error: 3989 SQLSTATE: HY000 (ER_TZ_TRANSITION_TYPE_TABLE_LOAD_ERROR)

  Message: Error while loading time zone description from mysql.time_zone_transition_type table

  ER_TZ_TRANSITION_TYPE_TABLE_LOAD_ERROR was added in 8.0.2.

- Error: 3990 SQLSTATE: HY000 (ER_TZ_TRANSITION_TABLE_TOO_MANY_TRANSITIONS)

  Message: Error while loading time zone description from mysql.time_zone_transition table: too much transitions

  ER_TZ_TRANSITION_TABLE_TOO_MANY_TRANSITIONS was added in 8.0.2.

- Error: 3991 SQLSTATE: HY000 (ER_TZ_TRANSITION_TABLE_BAD_TRANSITION_TYPE)

  Message: Error while loading time zone description from mysql.time_zone_transition table: bad transition type id

  ER_TZ_TRANSITION_TABLE_BAD_TRANSITION_TYPE was added in 8.0.2.

- Error: 3992 SQLSTATE: HY000 (ER_TZ_TRANSITION_TABLE_LOAD_ERROR)

  Message: Error while loading time zone description from mysql.time_zone_transition table

  ER_TZ_TRANSITION_TABLE_LOAD_ERROR was added in 8.0.2.

- Error: 3993 SQLSTATE: HY000 (ER_TZ_NO_TRANSITION_TYPES_IN_TIME_ZONE)

  Message: loading time zone without transition types

  ER_TZ_NO_TRANSITION_TYPES_IN_TIME_ZONE was added in 8.0.2.

- Error: 3994 SQLSTATE: HY000 (ER_TZ_OOM_LOADING_TIME_ZONE_DESCRIPTION)

  Message: Out of memory while loading time zone description

  ER_TZ_OOM_LOADING_TIME_ZONE_DESCRIPTION was added in 8.0.2.

- Error: 3995 SQLSTATE: HY000 (ER_TZ_CANT_BUILD_MKTIME_MAP)

  Message: Unable to build mktime map for time zone

  ER_TZ_CANT_BUILD_MKTIME_MAP was added in 8.0.2.

- Error: 3996 SQLSTATE: HY000 (ER_TZ_OOM_WHILE_LOADING_TIME_ZONE)

  Message: Out of memory while loading time zone

  ER_TZ_OOM_WHILE_LOADING_TIME_ZONE was added in 8.0.2.

- Error: 3997 SQLSTATE: HY000 (ER_TZ_OOM_WHILE_SETTING_TIME_ZONE)

  Message: Fatal error: Out of memory while setting new time zone

  ER_TZ_OOM_WHILE_SETTING_TIME_ZONE was added in 8.0.2.

- Error: 3998 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_UNTIL_CONDITION_BAD)

  Message: Slave SQL thread is stopped because UNTIL condition is bad(%s:%llu).

  ER_SLAVE_SQL_THREAD_STOPPED_UNTIL_CONDITION_BAD was added in 8.0.2.

- Error: 3999 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_UNTIL_POSITION_REACHED)

  Message: Slave SQL thread stopped because it reached its UNTIL position %llu

  ER_SLAVE_SQL_THREAD_STOPPED_UNTIL_POSITION_REACHED was added in 8.0.2.

- Error: 4000 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_BEFORE_GTIDS_ALREADY_APPLIED)

  Message: Slave SQL thread stopped because UNTIL SQL_BEFORE_GTIDS %s is already applied

  ER_SLAVE_SQL_THREAD_STOPPED_BEFORE_GTIDS_ALREADY_APPLIED was added in 8.0.2.

- Error: 4001 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_BEFORE_GTIDS_REACHED)

  Message: Slave SQL thread stopped because it reached UNTIL SQL_BEFORE_GTIDS %s

  ER_SLAVE_SQL_THREAD_STOPPED_BEFORE_GTIDS_REACHED was added in 8.0.2.

- Error: 4002 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_AFTER_GTIDS_REACHED)

  Message: Slave SQL thread stopped because it reached UNTIL SQL_AFTER_GTIDS %s

  ER_SLAVE_SQL_THREAD_STOPPED_AFTER_GTIDS_REACHED was added in 8.0.2.

- Error: 4003 SQLSTATE: HY000 (ER_SLAVE_SQL_THREAD_STOPPED_GAP_TRX_PROCESSED)

  Message: Slave SQL thread stopped according to UNTIL SQL_AFTER_MTS_GAPS as it has processed all gap transactions left from the previous slave session.

  ER_SLAVE_SQL_THREAD_STOPPED_GAP_TRX_PROCESSED was added in 8.0.2.

- Error: 4004 SQLSTATE: HY000 (ER_GROUP_REPLICATION_PLUGIN_NOT_INSTALLED)

  Message: Group Replication plugin is not installed.

  ER_GROUP_REPLICATION_PLUGIN_NOT_INSTALLED was added in 8.0.2.

- Error: 4005 SQLSTATE: HY000 (ER_GTID_ALREADY_ADDED_BY_USER)

  Message: The transaction owned GTID is already in the %s table, which is caused by an explicit modifying from user client.

  ER_GTID_ALREADY_ADDED_BY_USER was added in 8.0.2.

- Error: 4006 SQLSTATE: HY000 (ER_FAILED_TO_DELETE_FROM_GTID_EXECUTED_TABLE)

  Message: Failed to delete the row: '%s' from the gtid_executed table.

  ER_FAILED_TO_DELETE_FROM_GTID_EXECUTED_TABLE was added in 8.0.2.

- Error: 4007 SQLSTATE: HY000 (ER_FAILED_TO_COMPRESS_GTID_EXECUTED_TABLE)

  Message: Failed to compress the gtid_executed table.

  ER_FAILED_TO_COMPRESS_GTID_EXECUTED_TABLE was added in 8.0.2.

- Error: 4008 SQLSTATE: HY000 (ER_FAILED_TO_COMPRESS_GTID_EXECUTED_TABLE_OOM)

  Message: Failed to compress the gtid_executed table, because it is failed to allocate the THD.

  ER_FAILED_TO_COMPRESS_GTID_EXECUTED_TABLE_OOM was added in 8.0.2.

- Error: 4009 SQLSTATE: HY000 (ER_FAILED_TO_INIT_THREAD_ATTR_FOR_GTID_TABLE_COMPRESSION)

  Message: Failed to initialize thread attribute when creating compression thread.

  ER_FAILED_TO_INIT_THREAD_ATTR_FOR_GTID_TABLE_COMPRESSION was added in 8.0.2.

- Error: 4010 SQLSTATE: HY000 (ER_FAILED_TO_CREATE_GTID_TABLE_COMPRESSION_THREAD)

  Message: Can not create thread to compress gtid_executed table (errno= %d)

  ER_FAILED_TO_CREATE_GTID_TABLE_COMPRESSION_THREAD was added in 8.0.2.

- Error: 4011 SQLSTATE: HY000 (ER_FAILED_TO_JOIN_GTID_TABLE_COMPRESSION_THREAD)

  Message: Could not join gtid_executed table compression thread. error:%d

  ER_FAILED_TO_JOIN_GTID_TABLE_COMPRESSION_THREAD was added in 8.0.2.

- Error: 4012 SQLSTATE: HY000 (ER_NPIPE_FAILED_TO_INIT_SECURITY_DESCRIPTOR)

  Message: Can't start server : Initialize security descriptor: %s

  ER_NPIPE_FAILED_TO_INIT_SECURITY_DESCRIPTOR was added in 8.0.2.

- Error: 4013 SQLSTATE: HY000 (ER_NPIPE_FAILED_TO_SET_SECURITY_DESCRIPTOR)

  Message: Can't start server : Set security descriptor: %s

  ER_NPIPE_FAILED_TO_SET_SECURITY_DESCRIPTOR was added in 8.0.2.

- Error: 4014 SQLSTATE: HY000 (ER_NPIPE_PIPE_ALREADY_IN_USE)

  Message: Can't start server : Named Pipe "%s" already in use.

  ER_NPIPE_PIPE_ALREADY_IN_USE was added in 8.0.2.

- Error: 4015 SQLSTATE: HY000 (ER_NDB_SLAVE_SAW_EPOCH_LOWER_THAN_PREVIOUS_ON_START)

  Message: NDB Slave : At SQL thread start applying epoch %llu/%llu (%llu) from Master ServerId %u which is lower than previously applied epoch %llu/%llu (%llu). Group Master Log : %s Group Master Log Pos : %llu. Check slave positioning.

  ER_NDB_SLAVE_SAW_EPOCH_LOWER_THAN_PREVIOUS_ON_START was added in 8.0.2.

- Error: 4016 SQLSTATE: HY000 (ER_NDB_SLAVE_SAW_EPOCH_LOWER_THAN_PREVIOUS)

  Message: NDB Slave : SQL thread stopped as applying epoch %llu/%llu (%llu) from Master ServerId %u which is lower than previously applied epoch %llu/%llu (%llu). Group Master Log : %s Group Master Log Pos : %llu

  ER_NDB_SLAVE_SAW_EPOCH_LOWER_THAN_PREVIOUS was added in 8.0.2.

- Error: 4017 SQLSTATE: HY000 (ER_NDB_SLAVE_SAW_ALREADY_COMMITTED_EPOCH)

Message: NDB Slave : SQL thread stopped as attempted to reapply already committed epoch %llu/%llu (%llu) from server id %u. Group Master Log : %s Group Master Log Pos : %llu.

`ER_NDB_SLAVE_SAW_ALREADY_COMMITTED_EPOCH` was added in 8.0.2.

- Error: `4018` SQLSTATE: `HY000` (`ER_NDB_SLAVE_PREVIOUS_EPOCH_NOT_COMMITTED`)

  Message: NDB Slave : SQL thread stopped as attempting to apply new epoch %llu/%llu (%llu) while lower received epoch %llu/%llu (%llu) has not been committed. Master server id : %u. Group Master Log : %s Group Master Log Pos : %llu.

  `ER_NDB_SLAVE_PREVIOUS_EPOCH_NOT_COMMITTED` was added in 8.0.2.

- Error: `4019` SQLSTATE: `HY000` (`ER_NDB_SLAVE_MISSING_DATA_FOR_TIMESTAMP_COLUMN`)

  Message: NDB Slave: missing data for %s timestamp column %u.

  `ER_NDB_SLAVE_MISSING_DATA_FOR_TIMESTAMP_COLUMN` was added in 8.0.2.

- Error: `4020` SQLSTATE: `HY000` (`ER_NDB_SLAVE_LOGGING_EXCEPTIONS_TO`)

  Message: NDB Slave: Table %s.%s logging exceptions to %s.%s

  `ER_NDB_SLAVE_LOGGING_EXCEPTIONS_TO` was added in 8.0.2.

- Error: `4021` SQLSTATE: `HY000` (`ER_NDB_SLAVE_LOW_EPOCH_RESOLUTION`)

  Message: NDB Slave: Table %s.%s : %s, low epoch resolution

  `ER_NDB_SLAVE_LOW_EPOCH_RESOLUTION` was added in 8.0.2.

- Error: `4022` SQLSTATE: `HY000` (`ER_NDB_INFO_FOUND_UNEXPECTED_FIELD_TYPE`)

  Message: Found unexpected field type %u

  `ER_NDB_INFO_FOUND_UNEXPECTED_FIELD_TYPE` was added in 8.0.2.

- Error: `4023` SQLSTATE: `HY000` (`ER_NDB_INFO_FAILED_TO_CREATE_NDBINFO`)

  Message: Failed to create NdbInfo

  `ER_NDB_INFO_FAILED_TO_CREATE_NDBINFO` was added in 8.0.2.

- Error: `4024` SQLSTATE: `HY000` (`ER_NDB_INFO_FAILED_TO_INIT_NDBINFO`)

  Message: Failed to init NdbInfo

  `ER_NDB_INFO_FAILED_TO_INIT_NDBINFO` was added in 8.0.2.

- Error: `4025` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_WRONG_NUMBER_OF_FUNCTION_ARGUMENTS`)

  Message: ndb_serialize_cond: Unexpected mismatch of found and expected number of function arguments %u

  `ER_NDB_CLUSTER_WRONG_NUMBER_OF_FUNCTION_ARGUMENTS` was added in 8.0.2.

- Error: `4026` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_SCHEMA_INFO`)

  Message: %s - %s.%s

  `ER_NDB_CLUSTER_SCHEMA_INFO` was added in 8.0.2.

- Error: 4027 SQLSTATE: HY000 (ER_NDB_CLUSTER_GENERIC_MESSAGE)

  Message: %s

  ER_NDB_CLUSTER_GENERIC_MESSAGE was added in 8.0.2.

- Error: 4028 SQLSTATE: HY000 (ER_RPL_CANT_OPEN_INFO_TABLE)

  Message: Info table is not ready to be used. Table '%s.%s' cannot be opened.

  ER_RPL_CANT_OPEN_INFO_TABLE was added in 8.0.2.

- Error: 4029 SQLSTATE: HY000 (ER_RPL_CANT_SCAN_INFO_TABLE)

  Message: Info table is not ready to be used. Table '%s.%s' cannot be scanned.

  ER_RPL_CANT_SCAN_INFO_TABLE was added in 8.0.2.

- Error: 4030 SQLSTATE: HY000 (ER_RPL_CORRUPTED_INFO_TABLE)

  Message: Corrupted table %s.%s. Check out table definition.

  ER_RPL_CORRUPTED_INFO_TABLE was added in 8.0.2.

- Error: 4031 SQLSTATE: HY000 (ER_RPL_CORRUPTED_KEYS_IN_INFO_TABLE)

  Message: Info table has a problem with its key field(s). Table '%s.%s' expected field #%u to be '%s' but found '%s' instead.

  ER_RPL_CORRUPTED_KEYS_IN_INFO_TABLE was added in 8.0.2.

- Error: 4032 SQLSTATE: HY000 (ER_RPL_WORKER_ID_IS)

  Message: Choosing worker id %lu, the following is going to be %lu

  ER_RPL_WORKER_ID_IS was added in 8.0.2.

- Error: 4033 SQLSTATE: HY000 (ER_RPL_INCONSISTENT_TIMESTAMPS_IN_TRX)

  Message: Transaction is tagged with inconsistent logical timestamps: sequence_number (%lld) <= last_committed (%lld)

  ER_RPL_INCONSISTENT_TIMESTAMPS_IN_TRX was added in 8.0.2.

- Error: 4034 SQLSTATE: HY000 (ER_RPL_INCONSISTENT_SEQUENCE_NO_IN_TRX)

  Message: Transaction's sequence number is inconsistent with that of a preceding one: sequence_number (%lld) <= previous sequence_number (%lld)

  ER_RPL_INCONSISTENT_SEQUENCE_NO_IN_TRX was added in 8.0.2.

- Error: 4035 SQLSTATE: HY000 (ER_RPL_CHANNELS_REQUIRE_TABLES_AS_INFO_REPOSITORIES)

  Message: For the creation of replication channels the master info and relay log info repositories must be set to TABLE

  ER_RPL_CHANNELS_REQUIRE_TABLES_AS_INFO_REPOSITORIES was added in 8.0.2.

- Error: 4036 SQLSTATE: HY000 (ER_RPL_CHANNELS_REQUIRE_NON_ZERO_SERVER_ID)

  Message: For the creation of replication channels the server id must be different from 0

  ER_RPL_CHANNELS_REQUIRE_NON_ZERO_SERVER_ID was added in 8.0.2.

- Error: 4037 SQLSTATE: HY000 (ER_RPL_REPO_SHOULD_BE_TABLE)

  Message: Slave: Wrong repository. Repository should be TABLE

  ER_RPL_REPO_SHOULD_BE_TABLE was added in 8.0.2.

- Error: 4038 SQLSTATE: HY000 (ER_RPL_ERROR_CREATING_MASTER_INFO)

  Message: Error creating master info: %s.

  ER_RPL_ERROR_CREATING_MASTER_INFO was added in 8.0.2.

- Error: 4039 SQLSTATE: HY000 (ER_RPL_ERROR_CHANGING_MASTER_INFO_REPO_TYPE)

  Message: Error changing the type of master info's repository: %s.

  ER_RPL_ERROR_CHANGING_MASTER_INFO_REPO_TYPE was added in 8.0.2.

- Error: 4040 SQLSTATE: HY000 (ER_RPL_CHANGING_RELAY_LOG_INFO_REPO_TYPE_FAILED_DUE_TO_GAPS)

  Message: It is not possible to change the type of the relay log repository because there are workers repositories with possible execution gaps. The value of --relay_log_info_repository is altered to one of the found Worker repositories. The gaps have to be sorted out before resuming with the type change.

  ER_RPL_CHANGING_RELAY_LOG_INFO_REPO_TYPE_FAILED_DUE_TO_GAPS was added in 8.0.2.

- Error: 4041 SQLSTATE: HY000 (ER_RPL_ERROR_CREATING_RELAY_LOG_INFO)

  Message: Error creating relay log info: %s.

  ER_RPL_ERROR_CREATING_RELAY_LOG_INFO was added in 8.0.2.

- Error: 4042 SQLSTATE: HY000 (ER_RPL_ERROR_CHANGING_RELAY_LOG_INFO_REPO_TYPE)

  Message: Error changing the type of relay log info's repository: %s.

  ER_RPL_ERROR_CHANGING_RELAY_LOG_INFO_REPO_TYPE was added in 8.0.2.

- Error: 4043 SQLSTATE: HY000 (ER_RPL_FAILED_TO_DELETE_FROM_SLAVE_WORKERS_INFO_REPOSITORY)

  Message: Could not delete from Slave Workers info repository.

  ER_RPL_FAILED_TO_DELETE_FROM_SLAVE_WORKERS_INFO_REPOSITORY was added in 8.0.2.

- Error: 4044 SQLSTATE: HY000 (ER_RPL_FAILED_TO_RESET_STATE_IN_SLAVE_INFO_REPOSITORY)

  Message: Could not store the reset Slave Worker state into the slave info repository.

  ER_RPL_FAILED_TO_RESET_STATE_IN_SLAVE_INFO_REPOSITORY was added in 8.0.2.

- Error: 4045 SQLSTATE: HY000 (ER_RPL_ERROR_CHECKING_REPOSITORY)

  Message: Error in checking %s repository info type of %s.

  ER_RPL_ERROR_CHECKING_REPOSITORY was added in 8.0.2.

- Error: 4046 SQLSTATE: HY000 (ER_RPL_SLAVE_GENERIC_MESSAGE)

  Message: Slave: %s

ER_RPL_SLAVE_GENERIC_MESSAGE was added in 8.0.2.

- Error: 4047 SQLSTATE: HY000 (ER_RPL_SLAVE_COULD_NOT_CREATE_CHANNEL_LIST)

  Message: Slave: Could not create channel list

  ER_RPL_SLAVE_COULD_NOT_CREATE_CHANNEL_LIST was added in 8.0.2.

- Error: 4048 SQLSTATE: HY000 (ER_RPL_MULTISOURCE_REQUIRES_TABLE_TYPE_REPOSITORIES)

  Message: Slave: This slave was a multisourced slave previously which is supported only by both TABLE based master info and relay log info repositories. Found one or both of the info repos to be type FILE. Set both repos to type TABLE.

  ER_RPL_MULTISOURCE_REQUIRES_TABLE_TYPE_REPOSITORIES was added in 8.0.2.

- Error: 4049 SQLSTATE: HY000 (ER_RPL_SLAVE_FAILED_TO_INIT_A_MASTER_INFO_STRUCTURE)

  Message: Slave: Failed to initialize the master info structure for channel '%s'; its record may still be present in 'mysql.slave_master_info' table, consider deleting it.

  ER_RPL_SLAVE_FAILED_TO_INIT_A_MASTER_INFO_STRUCTURE was added in 8.0.2.

- Error: 4050 SQLSTATE: HY000 (ER_RPL_SLAVE_FAILED_TO_INIT_MASTER_INFO_STRUCTURE)

  Message: Failed to initialize the master info structure%s

  ER_RPL_SLAVE_FAILED_TO_INIT_MASTER_INFO_STRUCTURE was added in 8.0.2.

- Error: 4051 SQLSTATE: HY000 (ER_RPL_SLAVE_FAILED_TO_CREATE_CHANNEL_FROM_MASTER_INFO)

  Message: Slave: Failed to create a channel from master info table repository.

  ER_RPL_SLAVE_FAILED_TO_CREATE_CHANNEL_FROM_MASTER_INFO was added in 8.0.2.

- Error: 4052 SQLSTATE: HY000 (ER_RPL_FAILED_TO_CREATE_NEW_INFO_FILE)

  Message: Failed to create a new info file (file '%s', errno %d)

  ER_RPL_FAILED_TO_CREATE_NEW_INFO_FILE was added in 8.0.2.

- Error: 4053 SQLSTATE: HY000 (ER_RPL_FAILED_TO_CREATE_CACHE_FOR_INFO_FILE)

  Message: Failed to create a cache on info file (file '%s')

  ER_RPL_FAILED_TO_CREATE_CACHE_FOR_INFO_FILE was added in 8.0.2.

- Error: 4054 SQLSTATE: HY000 (ER_RPL_FAILED_TO_OPEN_INFO_FILE)

  Message: Failed to open the existing info file (file '%s', errno %d)

  ER_RPL_FAILED_TO_OPEN_INFO_FILE was added in 8.0.2.

- Error: 4055 SQLSTATE: HY000 (ER_RPL_GTID_MEMORY_FINALLY_AVAILABLE)

  Message: Server overcomes the temporary 'out of memory' in '%d' tries while allocating a new chunk of intervals for storing GTIDs.

  ER_RPL_GTID_MEMORY_FINALLY_AVAILABLE was added in 8.0.2.

- Error: 4056 SQLSTATE: HY000 (ER_SERVER_COST_UNKNOWN_COST_CONSTANT)

  Message: Unknown cost constant "%s" in mysql.server_cost table

  ER_SERVER_COST_UNKNOWN_COST_CONSTANT was added in 8.0.2.

- Error: 4057 SQLSTATE: HY000 (ER_SERVER_COST_INVALID_COST_CONSTANT)

  Message: Invalid value for cost constant "%s" in mysql.server_cost table: %.1f

  ER_SERVER_COST_INVALID_COST_CONSTANT was added in 8.0.2.

- Error: 4058 SQLSTATE: HY000 (ER_ENGINE_COST_UNKNOWN_COST_CONSTANT)

  Message: Unknown cost constant "%s" in mysql.engine_cost table

  ER_ENGINE_COST_UNKNOWN_COST_CONSTANT was added in 8.0.2.

- Error: 4059 SQLSTATE: HY000 (ER_ENGINE_COST_UNKNOWN_STORAGE_ENGINE)

  Message: Unknown storage engine "%s" in mysql.engine_cost table

  ER_ENGINE_COST_UNKNOWN_STORAGE_ENGINE was added in 8.0.2.

- Error: 4060 SQLSTATE: HY000 (ER_ENGINE_COST_INVALID_DEVICE_TYPE_FOR_SE)

  Message: Invalid device type %d for "%s" storage engine for cost constant "%s" in mysql.engine_cost table

  ER_ENGINE_COST_INVALID_DEVICE_TYPE_FOR_SE was added in 8.0.2.

- Error: 4061 SQLSTATE: HY000 (ER_ENGINE_COST_INVALID_CONST_CONSTANT_FOR_SE_AND_DEVICE)

  Message: Invalid value for cost constant "%s" for "%s" storage engine and device type %d in mysql.engine_cost table: %.1f

  ER_ENGINE_COST_INVALID_CONST_CONSTANT_FOR_SE_AND_DEVICE was added in 8.0.2.

- Error: 4062 SQLSTATE: HY000 (ER_SERVER_COST_FAILED_TO_READ)

  Message: init_read_record returned error when reading from mysql.server_cost table.

  ER_SERVER_COST_FAILED_TO_READ was added in 8.0.2.

- Error: 4063 SQLSTATE: HY000 (ER_ENGINE_COST_FAILED_TO_READ)

  Message: init_read_record returned error when reading from mysql.engine_cost table.

  ER_ENGINE_COST_FAILED_TO_READ was added in 8.0.2.

- Error: 4064 SQLSTATE: HY000 (ER_FAILED_TO_OPEN_COST_CONSTANT_TABLES)

  Message: Failed to open optimizer cost constant tables

  ER_FAILED_TO_OPEN_COST_CONSTANT_TABLES was added in 8.0.2.

- Error: 4065 SQLSTATE: HY000 (ER_RPL_UNSUPPORTED_UNIGNORABLE_EVENT_IN_STREAM)

  Message: Unsupported non-ignorable event fed into the event stream.

  ER_RPL_UNSUPPORTED_UNIGNORABLE_EVENT_IN_STREAM was added in 8.0.2.

- Error: 4066 SQLSTATE: HY000 (ER_RPL_GTID_LOG_EVENT_IN_STREAM)

Message: GTID_LOG_EVENT or ANONYMOUS_GTID_LOG_EVENT is not expected in an event stream %s.

`ER_RPL_GTID_LOG_EVENT_IN_STREAM` was added in 8.0.2.

- Error: `4067` SQLSTATE: `HY000` (`ER_RPL_UNEXPECTED_BEGIN_IN_STREAM`)

  Message: QUERY(BEGIN) is not expected in an event stream in the middle of a %s.

  `ER_RPL_UNEXPECTED_BEGIN_IN_STREAM` was added in 8.0.2.

- Error: `4068` SQLSTATE: `HY000` (`ER_RPL_UNEXPECTED_COMMIT_ROLLBACK_OR_XID_LOG_EVENT_IN_STREAM`)

  Message: QUERY(COMMIT or ROLLBACK) or XID_LOG_EVENT is not expected in an event stream %s.

  `ER_RPL_UNEXPECTED_COMMIT_ROLLBACK_OR_XID_LOG_EVENT_IN_STREAM` was added in 8.0.2.

- Error: `4069` SQLSTATE: `HY000` (`ER_RPL_UNEXPECTED_XA_ROLLBACK_IN_STREAM`)

  Message: QUERY(XA ROLLBACK) is not expected in an event stream %s.

  `ER_RPL_UNEXPECTED_XA_ROLLBACK_IN_STREAM` was added in 8.0.2.

- Error: `4070` SQLSTATE: `HY000` (`ER_EVENT_EXECUTION_FAILED_CANT_AUTHENTICATE_USER`)

  Message: Event Scheduler: [%s].[%s.%s] execution failed, failed to authenticate the user.

  `ER_EVENT_EXECUTION_FAILED_CANT_AUTHENTICATE_USER` was added in 8.0.2.

- Error: `4071` SQLSTATE: `HY000` (`ER_EVENT_EXECUTION_FAILED_USER_LOST_EVEN_PRIVILEGE`)

  Message: Event Scheduler: [%s].[%s.%s] execution failed, user no longer has EVENT privilege.

  `ER_EVENT_EXECUTION_FAILED_USER_LOST_EVEN_PRIVILEGE` was added in 8.0.2.

- Error: `4072` SQLSTATE: `HY000` (`ER_EVENT_ERROR_DURING_COMPILATION`)

  Message: Event Scheduler: %serror during compilation of %s.%s

  `ER_EVENT_ERROR_DURING_COMPILATION` was added in 8.0.2.

- Error: `4073` SQLSTATE: `HY000` (`ER_EVENT_DROPPING`)

  Message: Event Scheduler: Dropping %s.%s

  `ER_EVENT_DROPPING` was added in 8.0.2.

- Error: `4074` SQLSTATE: `HY000` (`ER_NDB_SCHEMA_GENERIC_MESSAGE`)

  Message: Ndb schema[%s.%s]: %s

  `ER_NDB_SCHEMA_GENERIC_MESSAGE` was added in 8.0.2.

- Error: `4075` SQLSTATE: `HY000` (`ER_RPL_INCOMPATIBLE_DECIMAL_IN_RBR`)

  Message: In RBR mode, Slave received incompatible DECIMAL field (old-style decimal field) from Master while creating conversion table. Please consider changing datatype on Master to new style decimal by executing ALTER command for column Name: %s.%s.%s.

ER_RPL_INCOMPATIBLE_DECIMAL_IN_RBR was added in 8.0.2.

- Error: 4076 SQLSTATE: HY000 (ER_INIT_ROOT_WITHOUT_PASSWORD)

  Message: root@localhost is created with an empty password ! Please consider switching off the --initialize-insecure option.

  ER_INIT_ROOT_WITHOUT_PASSWORD was added in 8.0.2.

- Error: 4077 SQLSTATE: HY000 (ER_INIT_GENERATING_TEMP_PASSWORD_FOR_ROOT)

  Message: A temporary password is generated for root@localhost: %s

  ER_INIT_GENERATING_TEMP_PASSWORD_FOR_ROOT was added in 8.0.2.

- Error: 4078 SQLSTATE: HY000 (ER_INIT_CANT_OPEN_BOOTSTRAP_FILE)

  Message: Failed to open the bootstrap file %s

  ER_INIT_CANT_OPEN_BOOTSTRAP_FILE was added in 8.0.2.

- Error: 4079 SQLSTATE: HY000 (ER_INIT_BOOTSTRAP_COMPLETE)

  Message: Bootstrapping complete

  ER_INIT_BOOTSTRAP_COMPLETE was added in 8.0.2.

- Error: 4080 SQLSTATE: HY000 (ER_INIT_DATADIR_NOT_EMPTY_WONT_INITIALIZE)

  Message: --initialize specified but the data directory has files in it. Aborting.

  ER_INIT_DATADIR_NOT_EMPTY_WONT_INITIALIZE was added in 8.0.2.

- Error: 4081 SQLSTATE: HY000 (ER_INIT_DATADIR_EXISTS_WONT_INITIALIZE)

  Message: --initialize specified on an existing data directory.

  ER_INIT_DATADIR_EXISTS_WONT_INITIALIZE was added in 8.0.2.

- Error: 4082 SQLSTATE: HY000 (ER_INIT_DATADIR_EXISTS_AND_PATH_TOO_LONG_WONT_INITIALIZE)

  Message: --initialize specified but the data directory exists and the path is too long. Aborting.

  ER_INIT_DATADIR_EXISTS_AND_PATH_TOO_LONG_WONT_INITIALIZE was added in 8.0.2.

- Error: 4083 SQLSTATE: HY000 (ER_INIT_DATADIR_EXISTS_AND_NOT_WRITABLE_WONT_INITIALIZE)

  Message: --initialize specified but the data directory exists and is not writable. Aborting.

  ER_INIT_DATADIR_EXISTS_AND_NOT_WRITABLE_WONT_INITIALIZE was added in 8.0.2.

- Error: 4084 SQLSTATE: HY000 (ER_INIT_CREATING_DD)

  Message: Creating the data directory %s

  ER_INIT_CREATING_DD was added in 8.0.2.

- Error: 4085 SQLSTATE: HY000 (ER_RPL_BINLOG_STARTING_DUMP)

  Message: Start binlog_dump to master_thread_id(%u) slave_server(%u), pos(%s, %llu)

ER_RPL_BINLOG_STARTING_DUMP was added in 8.0.2.

- Error: 4086 SQLSTATE: HY000 (ER_RPL_BINLOG_MASTER_SENDS_HEARTBEAT)

  Message: master sends heartbeat message

  ER_RPL_BINLOG_MASTER_SENDS_HEARTBEAT was added in 8.0.2.

- Error: 4087 SQLSTATE: HY000 (ER_RPL_BINLOG_SKIPPING_REMAINING_HEARTBEAT_INFO)

  Message: the rest of heartbeat info skipped ...

  ER_RPL_BINLOG_SKIPPING_REMAINING_HEARTBEAT_INFO was added in 8.0.2.

- Error: 4088 SQLSTATE: HY000 (ER_RPL_BINLOG_MASTER_USES_CHECKSUM_AND_SLAVE_CANT)

  Message: Master is configured to log replication events with checksum, but will not send such events to slaves that cannot process them

  ER_RPL_BINLOG_MASTER_USES_CHECKSUM_AND_SLAVE_CANT was added in 8.0.2.

- Error: 4089 SQLSTATE: HY000 (ER_NDB_QUERY_FAILED)

  Message: NDB: Query '%s' failed, error: %d: %s

  ER_NDB_QUERY_FAILED was added in 8.0.2.

- Error: 4090 SQLSTATE: HY000 (ER_KILLING_THREAD)

  Message: Killing thread %lu

  ER_KILLING_THREAD was added in 8.0.2.

- Error: 4091 SQLSTATE: HY000 (ER_DETACHING_SESSION_LEFT_BY_PLUGIN)

  Message: Plugin %s is deinitializing a thread but left a session attached. Detaching it forcefully.

  ER_DETACHING_SESSION_LEFT_BY_PLUGIN was added in 8.0.2.

- Error: 4092 SQLSTATE: HY000 (ER_CANT_DETACH_SESSION_LEFT_BY_PLUGIN)

  Message: Failed to detach the session.

  ER_CANT_DETACH_SESSION_LEFT_BY_PLUGIN was added in 8.0.2.

- Error: 4093 SQLSTATE: HY000 (ER_DETACHED_SESSIONS_LEFT_BY_PLUGIN)

  Message: Closed forcefully %u session%s left opened by plugin %s

  ER_DETACHED_SESSIONS_LEFT_BY_PLUGIN was added in 8.0.2.

- Error: 4094 SQLSTATE: HY000 (ER_FAILED_TO_DECREMENT_NUMBER_OF_THREADS)

  Message: Failed to decrement the number of threads

  ER_FAILED_TO_DECREMENT_NUMBER_OF_THREADS was added in 8.0.2.

- Error: 4095 SQLSTATE: HY000 (ER_PLUGIN_DID_NOT_DEINITIALIZE_THREADS)

  Message: Plugin %s did not deinitialize %u threads

  ER_PLUGIN_DID_NOT_DEINITIALIZE_THREADS was added in 8.0.2.

- Error: 4096 SQLSTATE: HY000 (ER_KILLED_THREADS_OF_PLUGIN)

  Message: Killed %u threads of plugin %s

  ER_KILLED_THREADS_OF_PLUGIN was added in 8.0.2.

- Error: 4097 SQLSTATE: HY000 (ER_NDB_SLAVE_MAX_REPLICATED_EPOCH_UNKNOWN)

  Message: NDB Slave : Could not determine maximum replicated epoch from %s.%s at Slave start, error %u %s

  ER_NDB_SLAVE_MAX_REPLICATED_EPOCH_UNKNOWN was added in 8.0.2.

- Error: 4098 SQLSTATE: HY000 (ER_NDB_SLAVE_MAX_REPLICATED_EPOCH_SET_TO)

  Message: NDB Slave : MaxReplicatedEpoch set to %llu (%u/%u) at Slave start

  ER_NDB_SLAVE_MAX_REPLICATED_EPOCH_SET_TO was added in 8.0.2.

- Error: 4099 SQLSTATE: HY000 (ER_NDB_NODE_ID_AND_MANAGEMENT_SERVER_INFO)

  Message: NDB: NodeID is %lu, management server '%s:%lu'

  ER_NDB_NODE_ID_AND_MANAGEMENT_SERVER_INFO was added in 8.0.2.

- Error: 4100 SQLSTATE: HY000 (ER_NDB_DISCONNECT_INFO)

  Message: tid %u: node[%u] transaction_hint=%u, transaction_no_hint=%u

  ER_NDB_DISCONNECT_INFO was added in 8.0.2.

- Error: 4101 SQLSTATE: HY000 (ER_NDB_COLUMN_DEFAULTS_DIFFER)

  Message: NDB Internal error: Default values differ for column %u, ndb_default: %d

  ER_NDB_COLUMN_DEFAULTS_DIFFER was added in 8.0.2.

- Error: 4102 SQLSTATE: HY000 (ER_NDB_COLUMN_SHOULD_NOT_HAVE_NATIVE_DEFAULT)

  Message: NDB Internal error: Column %u has native default, but shouldn't. Flags=%u, type=%u

  ER_NDB_COLUMN_SHOULD_NOT_HAVE_NATIVE_DEFAULT was added in 8.0.2.

- Error: 4103 SQLSTATE: HY000 (ER_NDB_FIELD_INFO)

  Message: field[name: '%s', type: %u, real_type: %u, flags: 0x%x, is_null: %d]

  ER_NDB_FIELD_INFO was added in 8.0.2.

- Error: 4104 SQLSTATE: HY000 (ER_NDB_COLUMN_INFO)

  Message: ndbCol[name: '%s', type: %u, column_no: %d, nullable: %d]

  ER_NDB_COLUMN_INFO was added in 8.0.2.

- Error: 4105 SQLSTATE: HY000 (ER_NDB_OOM_IN_FIX_UNIQUE_INDEX_ATTR_ORDER)

  Message: fix_unique_index_attr_order: my_malloc(%u) failure

  ER_NDB_OOM_IN_FIX_UNIQUE_INDEX_ATTR_ORDER was added in 8.0.2.

- Error: 4106 SQLSTATE: HY000 (ER_NDB_SLAVE_MALFORMED_EVENT_RECEIVED_ON_TABLE)

  Message: NDB Slave : Malformed event received on table %s cannot parse. Stopping Slave.

ER_NDB_SLAVE_MALFORMED_EVENT_RECEIVED_ON_TABLE was added in 8.0.2.

- Error: 4107 SQLSTATE: HY000 (ER_NDB_SLAVE_CONFLICT_FUNCTION_REQUIRES_ROLE)

  Message: NDB Slave : Conflict function %s defined on table %s requires ndb_slave_conflict_role variable to be set. Stopping slave.

  ER_NDB_SLAVE_CONFLICT_FUNCTION_REQUIRES_ROLE was added in 8.0.2.

- Error: 4108 SQLSTATE: HY000 (ER_NDB_SLAVE_CONFLICT_DETECTION_REQUIRES_TRANSACTION_IDS)

  Message: NDB Slave : Transactional conflict detection defined on table %s, but events received without transaction ids. Check --ndb-log-transaction-id setting on upstream Cluster.

  ER_NDB_SLAVE_CONFLICT_DETECTION_REQUIRES_TRANSACTION_IDS was added in 8.0.2.

- Error: 4109 SQLSTATE: HY000 (ER_NDB_SLAVE_BINLOG_MISSING_INFO_FOR_CONFLICT_DETECTION)

  Message: NDB Slave : Binlog event on table %s missing info necessary for conflict detection. Check binlog format options on upstream cluster.

  ER_NDB_SLAVE_BINLOG_MISSING_INFO_FOR_CONFLICT_DETECTION was added in 8.0.2.

- Error: 4110 SQLSTATE: HY000 (ER_NDB_ERROR_IN_READAUTOINCREMENTVALUE)

  Message: Error %lu in readAutoIncrementValue(): %s

  ER_NDB_ERROR_IN_READAUTOINCREMENTVALUE was added in 8.0.2.

- Error: 4111 SQLSTATE: HY000 (ER_NDB_FOUND_UNCOMMITTED_AUTOCOMMIT)

  Message: found uncommitted autocommit+rbwr transaction, commit status: %d

  ER_NDB_FOUND_UNCOMMITTED_AUTOCOMMIT was added in 8.0.2.

- Error: 4112 SQLSTATE: HY000 (ER_NDB_SLAVE_TOO_MANY_RETRIES)

  Message: Ndb slave retried transaction %u time(s) in vain. Giving up.

  ER_NDB_SLAVE_TOO_MANY_RETRIES was added in 8.0.2.

- Error: 4113 SQLSTATE: HY000 (ER_NDB_SLAVE_ERROR_IN_UPDATE_CREATE_INFO)

  Message: Error %lu in ::update_create_info(): %s

  ER_NDB_SLAVE_ERROR_IN_UPDATE_CREATE_INFO was added in 8.0.2.

- Error: 4114 SQLSTATE: HY000 (ER_NDB_SLAVE_CANT_ALLOCATE_TABLE_SHARE)

  Message: NDB: allocating table share for %s failed

  ER_NDB_SLAVE_CANT_ALLOCATE_TABLE_SHARE was added in 8.0.2.

- Error: 4115 SQLSTATE: HY000 (ER_NDB_BINLOG_ERROR_INFO_FROM_DA)

  Message: NDB Binlog: (%d)%s

  ER_NDB_BINLOG_ERROR_INFO_FROM_DA was added in 8.0.2.

- Error: 4116 SQLSTATE: HY000 (ER_NDB_BINLOG_CREATE_TABLE_EVENT)

Message: NDB Binlog: CREATE TABLE Event: %s

ER_NDB_BINLOG_CREATE_TABLE_EVENT was added in 8.0.2.

- Error: 4117 SQLSTATE: HY000 (ER_NDB_BINLOG_FAILED_CREATE_TABLE_EVENT_OPERATIONS)

  Message: NDB Binlog: FAILED CREATE TABLE event operations. Event: %s

  ER_NDB_BINLOG_FAILED_CREATE_TABLE_EVENT_OPERATIONS was added in 8.0.2.

- Error: 4118 SQLSTATE: HY000 (ER_NDB_BINLOG_RENAME_EVENT)

  Message: NDB Binlog: RENAME Event: %s

  ER_NDB_BINLOG_RENAME_EVENT was added in 8.0.2.

- Error: 4119 SQLSTATE: HY000 (ER_NDB_BINLOG_FAILED_CREATE_EVENT_OPERATIONS_DURING_RENAME)

  Message: NDB Binlog: FAILED create event operations during RENAME. Event %s

  ER_NDB_BINLOG_FAILED_CREATE_EVENT_OPERATIONS_DURING_RENAME was added in 8.0.2.

- Error: 4120 SQLSTATE: HY000 (ER_NDB_UNEXPECTED_RENAME_TYPE)

  Message: Unexpected rename case detected, sql_command: %d

  ER_NDB_UNEXPECTED_RENAME_TYPE was added in 8.0.2.

- Error: 4121 SQLSTATE: HY000 (ER_NDB_ERROR_IN_GET_AUTO_INCREMENT)

  Message: Error %lu in ::get_auto_increment(): %s

  ER_NDB_ERROR_IN_GET_AUTO_INCREMENT was added in 8.0.2.

- Error: 4122 SQLSTATE: HY000 (ER_NDB_CREATING_SHARE_IN_OPEN)

  Message: Calling ndbcluster_create_binlog_setup(%s) in ::open

  ER_NDB_CREATING_SHARE_IN_OPEN was added in 8.0.2.

- Error: 4123 SQLSTATE: HY000 (ER_NDB_TABLE_OPENED_READ_ONLY)

  Message: table '%s' opened read only

  ER_NDB_TABLE_OPENED_READ_ONLY was added in 8.0.2.

- Error: 4124 SQLSTATE: HY000 (ER_NDB_INITIALIZE_GIVEN_CLUSTER_PLUGIN_DISABLED)

  Message: NDB: '--initialize' -> ndbcluster plugin disabled

  ER_NDB_INITIALIZE_GIVEN_CLUSTER_PLUGIN_DISABLED was added in 8.0.2.

- Error: 4125 SQLSTATE: HY000 (ER_NDB_BINLOG_FORMAT_CHANGED_FROM_STMT_TO_MIXED)

  Message: NDB: Changed global value of binlog_format from STATEMENT to MIXED

  ER_NDB_BINLOG_FORMAT_CHANGED_FROM_STMT_TO_MIXED was added in 8.0.2.

- Error: 4126 SQLSTATE: HY000 (ER_NDB_TRAILING_SHARE_RELEASED_BY_CLOSE_CACHED_TABLES)

Message: NDB_SHARE: trailing share %s, released by close_cached_tables

`ER_NDB_TRAILING_SHARE_RELEASED_BY_CLOSE_CACHED_TABLES` was added in 8.0.2.

- Error: `4127` SQLSTATE: `HY000` (`ER_NDB_SHARE_ALREADY_EXISTS`)

  Message: NDB_SHARE: %s already exists use_count=%d. Moving away for safety, but possible memleak.

  `ER_NDB_SHARE_ALREADY_EXISTS` was added in 8.0.2.

- Error: `4128` SQLSTATE: `HY000` (`ER_NDB_HANDLE_TRAILING_SHARE_INFO`)

  Message: handle_trailing_share: %s use_count: %u

  `ER_NDB_HANDLE_TRAILING_SHARE_INFO` was added in 8.0.2.

- Error: `4129` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_GET_SHARE_INFO`)

  Message: ndbcluster_get_share: %s use_count: %u

  `ER_NDB_CLUSTER_GET_SHARE_INFO` was added in 8.0.2.

- Error: `4130` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_REAL_FREE_SHARE_INFO`)

  Message: ndbcluster_real_free_share: %s use_count: %u

  `ER_NDB_CLUSTER_REAL_FREE_SHARE_INFO` was added in 8.0.2.

- Error: `4131` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_REAL_FREE_SHARE_DROP_FAILED`)

  Message: ndbcluster_real_free_share: %s, still open - ignored 'free' (leaked?)

  `ER_NDB_CLUSTER_REAL_FREE_SHARE_DROP_FAILED` was added in 8.0.2.

- Error: `4132` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_FREE_SHARE_INFO`)

  Message: ndbcluster_free_share: %s use_count: %u

  `ER_NDB_CLUSTER_FREE_SHARE_INFO` was added in 8.0.2.

- Error: `4133` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_MARK_SHARE_DROPPED_INFO`)

  Message: ndbcluster_mark_share_dropped: %s use_count: %u

  `ER_NDB_CLUSTER_MARK_SHARE_DROPPED_INFO` was added in 8.0.2.

- Error: `4134` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_MARK_SHARE_DROPPED_DESTROYING_SHARE`)

  Message: ndbcluster_mark_share_dropped: destroys share %s

  `ER_NDB_CLUSTER_MARK_SHARE_DROPPED_DESTROYING_SHARE` was added in 8.0.2.

- Error: `4135` SQLSTATE: `HY000` (`ER_NDB_CLUSTER_OOM_THD_NDB`)

  Message: Could not allocate Thd_ndb object

  `ER_NDB_CLUSTER_OOM_THD_NDB` was added in 8.0.2.

- Error: `4136` SQLSTATE: `HY000` (`ER_NDB_BINLOG_NDB_TABLES_INITIALLY_READ_ONLY`)

  Message: NDB Binlog: Ndb tables initially read only.

ER_NDB_BINLOG_NDB_TABLES_INITIALLY_READ_ONLY was added in 8.0.2.

- Error: 4137 SQLSTATE: HY000 (ER_NDB_UTIL_THREAD_OOM)

  Message: ndb util thread: malloc failure, query cache not maintained properly

  ER_NDB_UTIL_THREAD_OOM was added in 8.0.2.

- Error: 4138 SQLSTATE: HY000 (ER_NDB_ILLEGAL_VALUE_FOR_NDB_RECV_THREAD_CPU_MASK)

  Message: Trying to set ndb_recv_thread_cpu_mask to illegal value = %s, ignored

  ER_NDB_ILLEGAL_VALUE_FOR_NDB_RECV_THREAD_CPU_MASK was added in 8.0.2.

- Error: 4139 SQLSTATE: HY000 (ER_NDB_TOO_MANY_CPUS_IN_NDB_RECV_THREAD_CPU_MASK)

  Message: Trying to set too many CPU's in ndb_recv_thread_cpu_mask, ignored this variable, erroneus value = %s

  ER_NDB_TOO_MANY_CPUS_IN_NDB_RECV_THREAD_CPU_MASK was added in 8.0.2.

- Error: 4140 SQLSTATE: HY000 (ER_DBUG_CHECK_SHARES_OPEN)

  Message: dbug_check_shares open:

  ER_DBUG_CHECK_SHARES_OPEN was added in 8.0.2.

- Error: 4141 SQLSTATE: HY000 (ER_DBUG_CHECK_SHARES_INFO)

  Message: %s.%s: state: %s(%u) use_count: %u

  ER_DBUG_CHECK_SHARES_INFO was added in 8.0.2.

- Error: 4142 SQLSTATE: HY000 (ER_DBUG_CHECK_SHARES_DROPPED)

  Message: dbug_check_shares dropped:

  ER_DBUG_CHECK_SHARES_DROPPED was added in 8.0.2.

- Error: 4143 SQLSTATE: HY000 (ER_INVALID_OR_OLD_TABLE_OR_DB_NAME)

  Message: Invalid (old?) table or database name '%s'

  ER_INVALID_OR_OLD_TABLE_OR_DB_NAME was added in 8.0.2.

- Error: 4144 SQLSTATE: HY000 (ER_TC_RECOVERING_AFTER_CRASH_USING)

  Message: Recovering after a crash using %s

  ER_TC_RECOVERING_AFTER_CRASH_USING was added in 8.0.2.

- Error: 4145 SQLSTATE: HY000 (ER_TC_CANT_AUTO_RECOVER_WITH_TC_HEURISTIC_RECOVER)

  Message: Cannot perform automatic crash recovery when --tc-heuristic-recover is used

  ER_TC_CANT_AUTO_RECOVER_WITH_TC_HEURISTIC_RECOVER was added in 8.0.2.

- Error: 4146 SQLSTATE: HY000 (ER_TC_BAD_MAGIC_IN_TC_LOG)

  Message: Bad magic header in tc log

ER_TC_BAD_MAGIC_IN_TC_LOG was added in 8.0.2.

- Error: 4147 SQLSTATE: HY000 (ER_TC_NEED_N_SE_SUPPORTING_2PC_FOR_RECOVERY)

  Message: Recovery failed! You must enable exactly %d storage engines that support two-phase commit protocol

  ER_TC_NEED_N_SE_SUPPORTING_2PC_FOR_RECOVERY was added in 8.0.2.

- Error: 4148 SQLSTATE: HY000 (ER_TC_RECOVERY_FAILED_THESE_ARE_YOUR_OPTIONS)

  Message: Crash recovery failed. Either correct the problem (if it's, for example, out of memory error) and restart, or delete tc log and start mysqld with --tc-heuristic-recover={commit|rollback}

  ER_TC_RECOVERY_FAILED_THESE_ARE_YOUR_OPTIONS was added in 8.0.2.

- Error: 4149 SQLSTATE: HY000 (ER_TC_HEURISTIC_RECOVERY_MODE)

  Message: Heuristic crash recovery mode

  ER_TC_HEURISTIC_RECOVERY_MODE was added in 8.0.2.

- Error: 4150 SQLSTATE: HY000 (ER_TC_HEURISTIC_RECOVERY_FAILED)

  Message: Heuristic crash recovery failed

  ER_TC_HEURISTIC_RECOVERY_FAILED was added in 8.0.2.

- Error: 4151 SQLSTATE: HY000 (ER_TC_RESTART_WITHOUT_TC_HEURISTIC_RECOVER)

  Message: Please restart mysqld without --tc-heuristic-recover

  ER_TC_RESTART_WITHOUT_TC_HEURISTIC_RECOVER was added in 8.0.2.

- Error: 4152 SQLSTATE: HY000 (ER_RPL_SLAVE_FAILED_TO_CREATE_OR_RECOVER_INFO_REPOSITORIES)

  Message: Failed to create or recover replication info repositories.

  ER_RPL_SLAVE_FAILED_TO_CREATE_OR_RECOVER_INFO_REPOSITORIES was added in 8.0.2.

- Error: 4153 SQLSTATE: HY000 (ER_RPL_SLAVE_AUTO_POSITION_IS_1_AND_GTID_MODE_IS_OFF)

  Message: Detected misconfiguration: replication channel '%s' was configured with AUTO_POSITION = 1, but the server was started with --gtid-mode=off. Either reconfigure replication using CHANGE MASTER TO MASTER_AUTO_POSITION = 0 FOR CHANNEL '%s', or change GTID_MODE to some value other than OFF, before starting the slave receiver thread.

  ER_RPL_SLAVE_AUTO_POSITION_IS_1_AND_GTID_MODE_IS_OFF was added in 8.0.2.

- Error: 4154 SQLSTATE: HY000 (ER_RPL_SLAVE_CANT_START_SLAVE_FOR_CHANNEL)

  Message: Slave: Could not start slave for channel '%s'. operation discontinued

  ER_RPL_SLAVE_CANT_START_SLAVE_FOR_CHANNEL was added in 8.0.2.

- Error: 4155 SQLSTATE: HY000 (ER_RPL_SLAVE_CANT_STOP_SLAVE_FOR_CHANNEL)

  Message: Slave: Could not stop slave for channel '%s' operation discontinued

  ER_RPL_SLAVE_CANT_STOP_SLAVE_FOR_CHANNEL was added in 8.0.2.

- Error: 4156 SQLSTATE: HY000 (ER_RPL_RECOVERY_NO_ROTATE_EVENT_FROM_MASTER)

  Message: Error during --relay-log-recovery: Could not locate rotate event from the master.

  ER_RPL_RECOVERY_NO_ROTATE_EVENT_FROM_MASTER was added in 8.0.2.

- Error: 4157 SQLSTATE: HY000 (ER_RPL_RECOVERY_ERROR_READ_RELAY_LOG)

  Message: Error during --relay-log-recovery: Error reading events from relay log: %d

  ER_RPL_RECOVERY_ERROR_READ_RELAY_LOG was added in 8.0.2.

- Error: 4158 SQLSTATE: HY000 (ER_RPL_RECOVERY_ERROR_FREEING_IO_CACHE)

  Message: Error during --relay-log-recovery: Error while freeing IO_CACHE object

  ER_RPL_RECOVERY_ERROR_FREEING_IO_CACHE was added in 8.0.2.

- Error: 4159 SQLSTATE: HY000 (ER_RPL_RECOVERY_SKIPPED_GROUP_REPLICATION_CHANNEL)

  Message: Relay log recovery skipped for group replication channel.

  ER_RPL_RECOVERY_SKIPPED_GROUP_REPLICATION_CHANNEL was added in 8.0.2.

- Error: 4160 SQLSTATE: HY000 (ER_RPL_RECOVERY_ERROR)

  Message: Error during --relay-log-recovery: %s

  ER_RPL_RECOVERY_ERROR was added in 8.0.2.

- Error: 4161 SQLSTATE: HY000 (ER_RPL_RECOVERY_IO_ERROR_READING_RELAY_LOG_INDEX)

  Message: Error during --relay-log-recovery: Could not read relay log index file due to an IO error.

  ER_RPL_RECOVERY_IO_ERROR_READING_RELAY_LOG_INDEX was added in 8.0.2.

- Error: 4162 SQLSTATE: HY000 (ER_RPL_RECOVERY_FILE_MASTER_POS_INFO)

  Message: Recovery from master pos %ld and file %s%s. Previous relay log pos and relay log file had been set to %lld, %s respectively.

  ER_RPL_RECOVERY_FILE_MASTER_POS_INFO was added in 8.0.2.

- Error: 4163 SQLSTATE: HY000 (ER_RPL_RECOVERY_REPLICATE_SAME_SERVER_ID_REQUIRES_POSITION)

  Message: Error during --relay-log-recovery: replicate_same_server_id is in use and sql thread's positions are not initialized, hence relay log recovery cannot happen.

  ER_RPL_RECOVERY_REPLICATE_SAME_SERVER_ID_REQUIRES_POSITION was added in 8.0.2.

- Error: 4164 SQLSTATE: HY000 (ER_RPL_MTS_RECOVERY_STARTING_COORDINATOR)

  Message: MTS recovery: starting coordinator thread to fill MTS gaps.

  ER_RPL_MTS_RECOVERY_STARTING_COORDINATOR was added in 8.0.2.

- Error: 4165 SQLSTATE: HY000 (ER_RPL_MTS_RECOVERY_FAILED_TO_START_COORDINATOR)

  Message: MTS recovery: failed to start the coordinator thread. Check the error log for additional details.

  ER_RPL_MTS_RECOVERY_FAILED_TO_START_COORDINATOR was added in 8.0.2.

- Error: 4166 SQLSTATE: HY000 (ER_RPL_MTS_AUTOMATIC_RECOVERY_FAILED)

  Message: MTS recovery: automatic recovery failed. Either the slave server had stopped due to an error during an earlier session or relay logs are corrupted.Fix the cause of the slave side error and restart the slave server or consider using RESET SLAVE.

  ER_RPL_MTS_AUTOMATIC_RECOVERY_FAILED was added in 8.0.2.

- Error: 4167 SQLSTATE: HY000 (ER_RPL_MTS_RECOVERY_CANT_OPEN_RELAY_LOG)

  Message: Failed to open the relay log '%s' (relay_log_pos %s).

  ER_RPL_MTS_RECOVERY_CANT_OPEN_RELAY_LOG was added in 8.0.2.

- Error: 4168 SQLSTATE: HY000 (ER_RPL_MTS_RECOVERY_SUCCESSFUL)

  Message: MTS recovery: completed successfully.

  ER_RPL_MTS_RECOVERY_SUCCESSFUL was added in 8.0.2.

- Error: 4169 SQLSTATE: HY000 (ER_RPL_SERVER_ID_MISSING)

  Message: Server id not set, will not start slave%s

  ER_RPL_SERVER_ID_MISSING was added in 8.0.2.

- Error: 4170 SQLSTATE: HY000 (ER_RPL_CANT_CREATE_SLAVE_THREAD)

  Message: Can't create slave thread%s.

  ER_RPL_CANT_CREATE_SLAVE_THREAD was added in 8.0.2.

- Error: 4171 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_WAS_KILLED)

  Message: The slave IO thread%s was killed while executing initialization query '%s'

  ER_RPL_SLAVE_IO_THREAD_WAS_KILLED was added in 8.0.2.

- Error: 4172 SQLSTATE: HY000 (ER_RPL_SLAVE_MASTER_UUID_HAS_CHANGED)

  Message: The master's UUID has changed, although this should not happen unless you have changed it manually. The old UUID was %s.

  ER_RPL_SLAVE_MASTER_UUID_HAS_CHANGED was added in 8.0.2.

- Error: 4173 SQLSTATE: HY000 (ER_RPL_SLAVE_USES_CHECKSUM_AND_MASTER_PRE_50)

  Message: Found a master with MySQL server version older than 5.0. With checksums enabled on the slave, replication might not work correctly. To ensure correct replication, restart the slave server with --slave_sql_verify_checksum=0.

  ER_RPL_SLAVE_USES_CHECKSUM_AND_MASTER_PRE_50 was added in 8.0.2.

- Error: 4174 SQLSTATE: HY000 (ER_RPL_SLAVE_SECONDS_BEHIND_MASTER_DUBIOUS)

  Message: "SELECT UNIX_TIMESTAMP()" failed on master, do not trust column Seconds_Behind_Master of SHOW SLAVE STATUS. Error: %s (%d)

  ER_RPL_SLAVE_SECONDS_BEHIND_MASTER_DUBIOUS was added in 8.0.2.

- Error: 4175 SQLSTATE: HY000 (ER_RPL_SLAVE_CANT_FLUSH_MASTER_INFO_FILE)

  Message: Failed to flush master info file.

ER_RPL_SLAVE_CANT_FLUSH_MASTER_INFO_FILE was added in 8.0.2.

- Error: 4176 SQLSTATE: HY000 (ER_RPL_SLAVE_REPORT_HOST_TOO_LONG)

  Message: The length of report_host is %zu. It is larger than the max length(%d), so this slave cannot be registered to the master%s.

  ER_RPL_SLAVE_REPORT_HOST_TOO_LONG was added in 8.0.2.

- Error: 4177 SQLSTATE: HY000 (ER_RPL_SLAVE_REPORT_USER_TOO_LONG)

  Message: The length of report_user is %zu. It is larger than the max length(%d), so this slave cannot be registered to the master%s.

  ER_RPL_SLAVE_REPORT_USER_TOO_LONG was added in 8.0.2.

- Error: 4178 SQLSTATE: HY000 (ER_RPL_SLAVE_REPORT_PASSWORD_TOO_LONG)

  Message: The length of report_password is %zu. It is larger than the max length(%d), so this slave cannot be registered to the master%s.

  ER_RPL_SLAVE_REPORT_PASSWORD_TOO_LONG was added in 8.0.2.

- Error: 4179 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_RETRYING)

  Message: Error on %s: %d %s, will retry in %d secs

  ER_RPL_SLAVE_ERROR_RETRYING was added in 8.0.2.

- Error: 4180 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_READING_FROM_SERVER)

  Message: Error reading packet from server%s: %s (server_errno=%d)

  ER_RPL_SLAVE_ERROR_READING_FROM_SERVER was added in 8.0.2.

- Error: 4181 SQLSTATE: HY000 (ER_RPL_SLAVE_DUMP_THREAD_KILLED_BY_MASTER)

  Message: Slave%s: received end packet from server due to dump thread being killed on master. Dump threads are killed for example during master shutdown, explicitly by a user, or when the master receives a binlog send request from a duplicate server UUID <%s> : Error %s

  ER_RPL_SLAVE_DUMP_THREAD_KILLED_BY_MASTER was added in 8.0.2.

- Error: 4182 SQLSTATE: HY000 (ER_RPL_MTS_STATISTICS)

  Message: Multi-threaded slave statistics%s: seconds elapsed = %lu; events assigned = %llu; worker queues filled over overrun level = %lu; waited due a Worker queue full = %lu; waited due the total size = %lu; waited at clock conflicts = %llu waited (count) when Workers occupied = %lu waited when Workers occupied = %llu

  ER_RPL_MTS_STATISTICS was added in 8.0.2.

- Error: 4183 SQLSTATE: HY000 (ER_RPL_MTS_RECOVERY_COMPLETE)

  Message: Slave%s: MTS Recovery has completed at relay log %s, position %llu master log %s, position %llu.

  ER_RPL_MTS_RECOVERY_COMPLETE was added in 8.0.2.

- Error: 4184 SQLSTATE: HY000 (ER_RPL_SLAVE_CANT_INIT_RELAY_LOG_POSITION)

  Message: Error initializing relay log position%s: %s

ER_RPL_SLAVE_CANT_INIT_RELAY_LOG_POSITION was added in 8.0.2.

- Error: 4185 SQLSTATE: HY000 (ER_RPL_SLAVE_CONNECTED_TO_MASTER_REPLICATION_STARTED)

  Message: Slave I/O thread%s: connected to master '%s@%s:%d',replication started in log '%s' at position %s

  ER_RPL_SLAVE_CONNECTED_TO_MASTER_REPLICATION_STARTED was added in 8.0.2.

- Error: 4186 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_KILLED)

  Message: Slave I/O thread%s killed while connecting to master

  ER_RPL_SLAVE_IO_THREAD_KILLED was added in 8.0.2.

- Error: 4187 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_CANT_REGISTER_ON_MASTER)

  Message: Slave I/O thread couldn't register on master

  ER_RPL_SLAVE_IO_THREAD_CANT_REGISTER_ON_MASTER was added in 8.0.2.

- Error: 4188 SQLSTATE: HY000 (ER_RPL_SLAVE_FORCING_TO_RECONNECT_IO_THREAD)

  Message: Forcing to reconnect slave I/O thread%s

  ER_RPL_SLAVE_FORCING_TO_RECONNECT_IO_THREAD was added in 8.0.2.

- Error: 4189 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_REQUESTING_BINLOG_DUMP)

  Message: Failed on request_dump()%s

  ER_RPL_SLAVE_ERROR_REQUESTING_BINLOG_DUMP was added in 8.0.2.

- Error: 4190 SQLSTATE: HY000 (ER_RPL_LOG_ENTRY_EXCEEDS_SLAVE_MAX_ALLOWED_PACKET)

  Message: Log entry on master is longer than slave_max_allowed_packet (%lu) on slave. If the entry is correct, restart the server with a higher value of slave_max_allowed_packet

  ER_RPL_LOG_ENTRY_EXCEEDS_SLAVE_MAX_ALLOWED_PACKET was added in 8.0.2.

- Error: 4191 SQLSTATE: HY000 (ER_RPL_SLAVE_STOPPING_AS_MASTER_OOM)

  Message: Stopping slave I/O thread due to out-of-memory error from master

  ER_RPL_SLAVE_STOPPING_AS_MASTER_OOM was added in 8.0.2.

- Error: 4192 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_ABORTED_WAITING_FOR_RELAY_LOG_SPACE)

  Message: Slave I/O thread aborted while waiting for relay log space

  ER_RPL_SLAVE_IO_THREAD_ABORTED_WAITING_FOR_RELAY_LOG_SPACE was added in 8.0.2.

- Error: 4193 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_EXITING)

  Message: Slave I/O thread exiting%s, read up to log '%s', position %s

  ER_RPL_SLAVE_IO_THREAD_EXITING was added in 8.0.2.

- Error: 4194 SQLSTATE: HY000 (ER_RPL_SLAVE_CANT_INITIALIZE_SLAVE_WORKER)

Message: Failed during slave worker initialization%s

ER_RPL_SLAVE_CANT_INITIALIZE_SLAVE_WORKER was added in 8.0.2.

- Error: 4195 SQLSTATE: HY000 (ER_RPL_MTS_GROUP_RECOVERY_RELAY_LOG_INFO_FOR_WORKER)

  Message: Slave: MTS group recovery relay log info based on Worker-Id %lu, group_relay_log_name %s, group_relay_log_pos %llu group_master_log_name %s, group_master_log_pos %llu

  ER_RPL_MTS_GROUP_RECOVERY_RELAY_LOG_INFO_FOR_WORKER was added in 8.0.2.

- Error: 4196 SQLSTATE: HY000 (ER_RPL_ERROR_LOOKING_FOR_LOG)

  Message: Error looking for %s.

  ER_RPL_ERROR_LOOKING_FOR_LOG was added in 8.0.2.

- Error: 4197 SQLSTATE: HY000 (ER_RPL_MTS_GROUP_RECOVERY_RELAY_LOG_INFO)

  Message: Slave: MTS group recovery relay log info group_master_log_name %s, event_master_log_pos %llu.

  ER_RPL_MTS_GROUP_RECOVERY_RELAY_LOG_INFO was added in 8.0.2.

- Error: 4198 SQLSTATE: HY000 (ER_RPL_CANT_FIND_FOLLOWUP_FILE)

  Message: Error looking for file after %s.

  ER_RPL_CANT_FIND_FOLLOWUP_FILE was added in 8.0.2.

- Error: 4199 SQLSTATE: HY000 (ER_RPL_MTS_CHECKPOINT_PERIOD_DIFFERS_FROM_CNT)

  Message: This an error cnt != mts_checkpoint_period

  ER_RPL_MTS_CHECKPOINT_PERIOD_DIFFERS_FROM_CNT was added in 8.0.2.

- Error: 4200 SQLSTATE: HY000 (ER_RPL_SLAVE_WORKER_THREAD_CREATION_FAILED)

  Message: Failed during slave worker thread creation%s

  ER_RPL_SLAVE_WORKER_THREAD_CREATION_FAILED was added in 8.0.2.

- Error: 4201 SQLSTATE: HY000 (ER_RPL_SLAVE_WORKER_THREAD_CREATION_FAILED_WITH_ERRNO)

  Message: Failed during slave worker thread creation%s (errno= %d)

  ER_RPL_SLAVE_WORKER_THREAD_CREATION_FAILED_WITH_ERRNO was added in 8.0.2.

- Error: 4202 SQLSTATE: HY000 (ER_RPL_SLAVE_FAILED_TO_INIT_PARTITIONS_HASH)

  Message: Failed to init partitions hash

  ER_RPL_SLAVE_FAILED_TO_INIT_PARTITIONS_HASH was added in 8.0.2.

- Error: 4203 SQLSTATE: HY000 (ER_RPL_SLAVE_NDB_TABLES_NOT_AVAILABLE)

  Message: Slave SQL thread : NDB : Tables not available after %lu seconds. Consider increasing --ndb-wait-setup value

  ER_RPL_SLAVE_NDB_TABLES_NOT_AVAILABLE was added in 8.0.2.

- Error: 4204 SQLSTATE: HY000 (ER_RPL_SLAVE_SQL_THREAD_STARTING)

  Message: Slave SQL thread%s initialized, starting replication in log '%s' at position %s, relay log '%s' position: %s

  ER_RPL_SLAVE_SQL_THREAD_STARTING was added in 8.0.2.

- Error: 4205 SQLSTATE: HY000 (ER_RPL_SLAVE_SKIP_COUNTER_EXECUTED)

  Message: 'SQL_SLAVE_SKIP_COUNTER=%ld' executed at relay_log_file='%s', relay_log_pos='%ld', master_log_name='%s', master_log_pos='%ld' and new position at relay_log_file='%s', relay_log_pos='%ld', master_log_name='%s', master_log_pos='%ld'

  ER_RPL_SLAVE_SKIP_COUNTER_EXECUTED was added in 8.0.2.

- Error: 4206 SQLSTATE: HY000 (ER_RPL_SLAVE_ADDITIONAL_ERROR_INFO_FROM_DA)

  Message: Slave (additional info): %s Error_code: %d

  ER_RPL_SLAVE_ADDITIONAL_ERROR_INFO_FROM_DA was added in 8.0.2.

- Error: 4207 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_INFO_FROM_DA)

  Message: Slave: %s Error_code: %d

  ER_RPL_SLAVE_ERROR_INFO_FROM_DA was added in 8.0.2.

- Error: 4208 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_LOADING_USER_DEFINED_LIBRARY)

  Message: Error loading user-defined library, slave SQL thread aborted. Install the missing library, and restart the slave SQL thread with "SLAVE START". We stopped at log '%s' position %s

  ER_RPL_SLAVE_ERROR_LOADING_USER_DEFINED_LIBRARY was added in 8.0.2.

- Error: 4209 SQLSTATE: HY000 (ER_RPL_SLAVE_ERROR_RUNNING_QUERY)

  Message: Error running query, slave SQL thread aborted. Fix the problem, and restart the slave SQL thread with "SLAVE START". We stopped at log '%s' position %s

  ER_RPL_SLAVE_ERROR_RUNNING_QUERY was added in 8.0.2.

- Error: 4210 SQLSTATE: HY000 (ER_RPL_SLAVE_SQL_THREAD_EXITING)

  Message: Slave SQL thread%s exiting, replication stopped in log '%s' at position %s

  ER_RPL_SLAVE_SQL_THREAD_EXITING was added in 8.0.2.

- Error: 4211 SQLSTATE: HY000 (ER_RPL_SLAVE_READ_INVALID_EVENT_FROM_MASTER)

  Message: Read invalid event from master: '%s', master could be corrupt but a more likely cause of this is a bug

  ER_RPL_SLAVE_READ_INVALID_EVENT_FROM_MASTER was added in 8.0.2.

- Error: 4212 SQLSTATE: HY000 (ER_RPL_SLAVE_QUEUE_EVENT_FAILED_INVALID_CONFIGURATION)

  Message: The queue event failed for channel '%s' as its configuration is invalid.

  ER_RPL_SLAVE_QUEUE_EVENT_FAILED_INVALID_CONFIGURATION was added in 8.0.2.

- Error: 4213 SQLSTATE: HY000 (ER_RPL_SLAVE_IO_THREAD_DETECTED_UNEXPECTED_EVENT_SEQUENCE)

Message: An unexpected event sequence was detected by the IO thread while queuing the event received from master '%s' binary log file, at position %llu.

`ER_RPL_SLAVE_IO_THREAD_DETECTED_UNEXPECTED_EVENT_SEQUENCE` was added in 8.0.2.

- Error: `4214` SQLSTATE: `HY000` (`ER_RPL_SLAVE_CANT_USE_CHARSET`)

  Message: '%s' can not be used as client character set. '%s' will be used as default client character set while connecting to master.

  `ER_RPL_SLAVE_CANT_USE_CHARSET` was added in 8.0.2.

- Error: `4215` SQLSTATE: `HY000` (`ER_RPL_SLAVE_CONNECTED_TO_MASTER_REPLICATION_RESUMED`)

  Message: Slave%s: connected to master '%s@%s:%d',replication resumed in log '%s' at position %s

  `ER_RPL_SLAVE_CONNECTED_TO_MASTER_REPLICATION_RESUMED` was added in 8.0.2.

- Error: `4216` SQLSTATE: `HY000` (`ER_RPL_SLAVE_NEXT_LOG_IS_ACTIVE`)

  Message: next log '%s' is active

  `ER_RPL_SLAVE_NEXT_LOG_IS_ACTIVE` was added in 8.0.2.

- Error: `4217` SQLSTATE: `HY000` (`ER_RPL_SLAVE_NEXT_LOG_IS_INACTIVE`)

  Message: next log '%s' is not active

  `ER_RPL_SLAVE_NEXT_LOG_IS_INACTIVE` was added in 8.0.2.

- Error: `4218` SQLSTATE: `HY000` (`ER_RPL_SLAVE_SQL_THREAD_IO_ERROR_READING_EVENT`)

  Message: Slave SQL thread%s: I/O error reading event (errno: %d cur_log->error: %d)

  `ER_RPL_SLAVE_SQL_THREAD_IO_ERROR_READING_EVENT` was added in 8.0.2.

- Error: `4219` SQLSTATE: `HY000` (`ER_RPL_SLAVE_ERROR_READING_RELAY_LOG_EVENTS`)

  Message: Error reading relay log event%s: %s

  `ER_RPL_SLAVE_ERROR_READING_RELAY_LOG_EVENTS` was added in 8.0.2.

- Error: `4220` SQLSTATE: `HY000` (`ER_SLAVE_CHANGE_MASTER_TO_EXECUTED`)

  Message: 'CHANGE MASTER TO%s executed'. Previous state master_host='%s', master_port= %u, master_log_file='%s', master_log_pos= %ld, master_bind='%s'. New state master_host='%s', master_port= %u, master_log_file='%s', master_log_pos= %ld, master_bind='%s'.

  `ER_SLAVE_CHANGE_MASTER_TO_EXECUTED` was added in 8.0.2.

- Error: `4221` SQLSTATE: `HY000` (`ER_RPL_SLAVE_NEW_MASTER_INFO_NEEDS_REPOS_TYPE_OTHER_THAN_FILE`)

  Message: Slave: Cannot create new master info structure when repositories are of type FILE. Convert slave repositories to TABLE to replicate from multiple sources.

  `ER_RPL_SLAVE_NEW_MASTER_INFO_NEEDS_REPOS_TYPE_OTHER_THAN_FILE` was added in 8.0.2.

- Error: `4222` SQLSTATE: `HY000` (`ER_RPL_FAILED_TO_STAT_LOG_IN_INDEX`)

  Message: log %s listed in the index, but failed to stat.

ER_RPL_FAILED_TO_STAT_LOG_IN_INDEX was added in 8.0.2.

- Error: 4223 SQLSTATE: HY000 (ER_RPL_LOG_NOT_FOUND_WHILE_COUNTING_RELAY_LOG_SPACE)

  Message: Could not find first log while counting relay log space.

  ER_RPL_LOG_NOT_FOUND_WHILE_COUNTING_RELAY_LOG_SPACE was added in 8.0.2.

- Error: 4224 SQLSTATE: HY000 (ER_SLAVE_CANT_USE_TEMPDIR)

  Message: Unable to use slave's temporary directory '%s'.

  ER_SLAVE_CANT_USE_TEMPDIR was added in 8.0.2.

- Error: 4225 SQLSTATE: HY000 (ER_RPL_RELAY_LOG_NEEDS_FILE_NOT_DIRECTORY)

  Message: Path '%s' is a directory name, please specify a file name for --relay-log option.

  ER_RPL_RELAY_LOG_NEEDS_FILE_NOT_DIRECTORY was added in 8.0.2.

- Error: 4226 SQLSTATE: HY000 (ER_RPL_RELAY_LOG_INDEX_NEEDS_FILE_NOT_DIRECTORY)

  Message: Path '%s' is a directory name, please specify a file name for --relay-log-index option.

  ER_RPL_RELAY_LOG_INDEX_NEEDS_FILE_NOT_DIRECTORY was added in 8.0.2.

- Error: 4227 SQLSTATE: HY000 (ER_RPL_PLEASE_USE_OPTION_RELAY_LOG)

  Message: Neither --relay-log nor --relay-log-index were used; so replication may break when this MySQL server acts as a slave and has his hostname changed!! Please use '--relay-log=%s' to avoid this problem.

  ER_RPL_PLEASE_USE_OPTION_RELAY_LOG was added in 8.0.2.

- Error: 4228 SQLSTATE: HY000 (ER_RPL_OPEN_INDEX_FILE_FAILED)

  Message: Failed in open_index_file() called from Relay_log_info::rli_init_info().

  ER_RPL_OPEN_INDEX_FILE_FAILED was added in 8.0.2.

- Error: 4229 SQLSTATE: HY000 (ER_RPL_CANT_INITIALIZE_GTID_SETS_IN_RLI_INIT_INFO)

  Message: Failed in init_gtid_sets() called from Relay_log_info::rli_init_info().

  ER_RPL_CANT_INITIALIZE_GTID_SETS_IN_RLI_INIT_INFO was added in 8.0.2.

- Error: 4230 SQLSTATE: HY000 (ER_RPL_CANT_OPEN_LOG_IN_RLI_INIT_INFO)

  Message: Failed in open_log() called from Relay_log_info::rli_init_info().

  ER_RPL_CANT_OPEN_LOG_IN_RLI_INIT_INFO was added in 8.0.2.

- Error: 4231 SQLSTATE: HY000 (ER_RPL_ERROR_WRITING_RELAY_LOG_CONFIGURATION)

  Message: Error writing relay log configuration.

  ER_RPL_ERROR_WRITING_RELAY_LOG_CONFIGURATION was added in 8.0.2.

- Error: 4232 SQLSTATE: HY000 (ER_NDB_OOM_GET_NDB_BLOBS_VALUE)

  Message: get_ndb_blobs_value: my_malloc(%u) failed

ER_NDB_OOM_GET_NDB_BLOBS_VALUE was added in 8.0.2.

- Error: 4233 SQLSTATE: HY000 (ER_NDB_THREAD_TIMED_OUT)

  Message: NDB: Thread id %u timed out (30s) waiting for epoch %u/%u to be handled. Progress : %u/%u -> %u/%u.

  ER_NDB_THREAD_TIMED_OUT was added in 8.0.2.

- Error: 4234 SQLSTATE: HY000 (ER_NDB_TABLE_IS_NOT_DISTRIBUTED)

  Message: NDB: Inconsistency detected in distributed privilege tables. Table '%s.%s' is not distributed

  ER_NDB_TABLE_IS_NOT_DISTRIBUTED was added in 8.0.2.

- Error: 4235 SQLSTATE: HY000 (ER_NDB_CREATING_TABLE)

  Message: NDB: Creating %s.%s

  ER_NDB_CREATING_TABLE was added in 8.0.2.

- Error: 4236 SQLSTATE: HY000 (ER_NDB_FLUSHING_TABLE_INFO)

  Message: NDB: Flushing %s.%s

  ER_NDB_FLUSHING_TABLE_INFO was added in 8.0.2.

- Error: 4237 SQLSTATE: HY000 (ER_NDB_CLEANING_STRAY_TABLES)

  Message: NDB: Cleaning stray tables from database '%s'

  ER_NDB_CLEANING_STRAY_TABLES was added in 8.0.2.

- Error: 4238 SQLSTATE: HY000 (ER_NDB_DISCOVERED_MISSING_DB)

  Message: NDB: Discovered missing database '%s'

  ER_NDB_DISCOVERED_MISSING_DB was added in 8.0.2.

- Error: 4239 SQLSTATE: HY000 (ER_NDB_DISCOVERED_REMAINING_DB)

  Message: NDB: Discovered remaining database '%s'

  ER_NDB_DISCOVERED_REMAINING_DB was added in 8.0.2.

- Error: 4240 SQLSTATE: HY000 (ER_NDB_CLUSTER_FIND_ALL_DBS_RETRY)

  Message: NDB: ndbcluster_find_all_databases retry: %u - %s

  ER_NDB_CLUSTER_FIND_ALL_DBS_RETRY was added in 8.0.2.

- Error: 4241 SQLSTATE: HY000 (ER_NDB_CLUSTER_FIND_ALL_DBS_FAIL)

  Message: NDB: ndbcluster_find_all_databases fail: %u - %s

  ER_NDB_CLUSTER_FIND_ALL_DBS_FAIL was added in 8.0.2.

- Error: 4242 SQLSTATE: HY000 (ER_NDB_SKIPPING_SETUP_TABLE)

  Message: NDB: skipping setup table %s.%s, in state %d

  ER_NDB_SKIPPING_SETUP_TABLE was added in 8.0.2.

- Error: 4243 SQLSTATE: HY000 (ER_NDB_FAILED_TO_SET_UP_TABLE)

Message: NDB: failed to setup table %s.%s, error: %d, %s

ER_NDB_FAILED_TO_SET_UP_TABLE was added in 8.0.2.

- Error: 4244 SQLSTATE: HY000 (ER_NDB_MISSING_FRM_DISCOVERING)

  Message: NDB: missing frm for %s.%s, discovering...

  ER_NDB_MISSING_FRM_DISCOVERING was added in 8.0.2.

- Error: 4245 SQLSTATE: HY000 (ER_NDB_MISMATCH_IN_FRM_DISCOVERING)

  Message: NDB: mismatch in frm for %s.%s, discovering...

  ER_NDB_MISMATCH_IN_FRM_DISCOVERING was added in 8.0.2.

- Error: 4246 SQLSTATE: HY000 (ER_NDB_BINLOG_CLEANING_UP_SETUP_LEFTOVERS)

  Message: ndb_binlog_setup: Clean up leftovers

  ER_NDB_BINLOG_CLEANING_UP_SETUP_LEFTOVERS was added in 8.0.2.

- Error: 4247 SQLSTATE: HY000 (ER_NDB_WAITING_INFO)

  Message: NDB %s: waiting max %u sec for %s %s. epochs: (%u/%u,%u/%u,%u/%u) injector proc_info: %s

  ER_NDB_WAITING_INFO was added in 8.0.2.

- Error: 4248 SQLSTATE: HY000 (ER_NDB_WAITING_INFO_WITH_MAP)

  Message: NDB %s: waiting max %u sec for %s %s. epochs: (%u/%u,%u/%u,%u/%u) injector proc_info: %s map: %x%08x

  ER_NDB_WAITING_INFO_WITH_MAP was added in 8.0.2.

- Error: 4249 SQLSTATE: HY000 (ER_NDB_TIMEOUT_WHILE_DISTRIBUTING)

  Message: NDB %s: distributing %s timed out. Ignoring...

  ER_NDB_TIMEOUT_WHILE_DISTRIBUTING was added in 8.0.2.

- Error: 4250 SQLSTATE: HY000 (ER_NDB_NOT_WAITING_FOR_DISTRIBUTING)

  Message: NDB %s: not waiting for distributing %s

  ER_NDB_NOT_WAITING_FOR_DISTRIBUTING was added in 8.0.2.

- Error: 4251 SQLSTATE: HY000 (ER_NDB_DISTRIBUTED_INFO)

  Message: NDB: distributed %s.%s(%u/%u) type: %s(%u) query: \'%s\' to %x%08x

  ER_NDB_DISTRIBUTED_INFO was added in 8.0.2.

- Error: 4252 SQLSTATE: HY000 (ER_NDB_DISTRIBUTION_COMPLETE)

  Message: NDB: distribution of %s.%s(%u/%u) type: %s(%u) query: \'%s\' - complete!

  ER_NDB_DISTRIBUTION_COMPLETE was added in 8.0.2.

- Error: 4253 SQLSTATE: HY000 (ER_NDB_SCHEMA_DISTRIBUTION_FAILED)

  Message: NDB Schema dist: Data node: %d failed, subscriber bitmask %x%08x

ER_NDB_SCHEMA_DISTRIBUTION_FAILED was added in 8.0.2.

- Error: 4254 SQLSTATE: HY000 (ER_NDB_SCHEMA_DISTRIBUTION_REPORTS_SUBSCRIBE)

  Message: NDB Schema dist: Data node: %d reports subscribe from node %d, subscriber bitmask %x %08x

  ER_NDB_SCHEMA_DISTRIBUTION_REPORTS_SUBSCRIBE was added in 8.0.2.

- Error: 4255 SQLSTATE: HY000 (ER_NDB_SCHEMA_DISTRIBUTION_REPORTS_UNSUBSCRIBE)

  Message: NDB Schema dist: Data node: %d reports unsubscribe from node %d, subscriber bitmask %x%08x

  ER_NDB_SCHEMA_DISTRIBUTION_REPORTS_UNSUBSCRIBE was added in 8.0.2.

- Error: 4256 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_DISCOVER_TABLE_FROM_SCHEMA_EVENT)

  Message: NDB Binlog: Could not discover table '%s.%s' from binlog schema event '%s' from node %d. my_errno: %d

  ER_NDB_BINLOG_CANT_DISCOVER_TABLE_FROM_SCHEMA_EVENT was added in 8.0.2.

- Error: 4257 SQLSTATE: HY000 (ER_NDB_BINLOG_SIGNALLING_UNKNOWN_VALUE)

  Message: NDB: unknown value for binlog signalling 0x%X, %s not logged

  ER_NDB_BINLOG_SIGNALLING_UNKNOWN_VALUE was added in 8.0.2.

- Error: 4258 SQLSTATE: HY000 (ER_NDB_BINLOG_REPLY_TO)

  Message: NDB: reply to %s.%s(%u/%u) from %s to %x%08x

  ER_NDB_BINLOG_REPLY_TO was added in 8.0.2.

- Error: 4259 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_RELEASE_SLOCK)

  Message: NDB: Could not release slock on '%s.%s', Error code: %d Message: %s

  ER_NDB_BINLOG_CANT_RELEASE_SLOCK was added in 8.0.2.

- Error: 4260 SQLSTATE: HY000 (ER_NDB_CANT_FIND_TABLE)

  Message: NDB schema: Could not find table '%s.%s' in NDB

  ER_NDB_CANT_FIND_TABLE was added in 8.0.2.

- Error: 4261 SQLSTATE: HY000 (ER_NDB_DISCARDING_EVENT_NO_OBJ)

  Message: NDB: Discarding event...no obj: %s (%u/%u)

  ER_NDB_DISCARDING_EVENT_NO_OBJ was added in 8.0.2.

- Error: 4262 SQLSTATE: HY000 (ER_NDB_DISCARDING_EVENT_ID_VERSION_MISMATCH)

  Message: NDB: Discarding event...key: %s non matching id/version [%u/%u] != [%u/%u]

  ER_NDB_DISCARDING_EVENT_ID_VERSION_MISMATCH was added in 8.0.2.

- Error: 4263 SQLSTATE: HY000 (ER_NDB_CLEAR_SLOCK_INFO)

  Message: NDB: CLEAR_SLOCK key: %s(%u/%u) %x%08x, from %s to %x%08x

ER_NDB_CLEAR_SLOCK_INFO was added in 8.0.2.

- Error: 4264 SQLSTATE: HY000 (ER_NDB_BINLOG_SKIPPING_LOCAL_TABLE)

  Message: NDB Binlog: Skipping locally defined table '%s.%s' from binlog schema event '%s' from node %d.

  ER_NDB_BINLOG_SKIPPING_LOCAL_TABLE was added in 8.0.2.

- Error: 4265 SQLSTATE: HY000 (ER_NDB_BINLOG_ONLINE_ALTER_RENAME)

  Message: NDB Binlog: handling online alter/rename

  ER_NDB_BINLOG_ONLINE_ALTER_RENAME was added in 8.0.2.

- Error: 4266 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_REOPEN_SHADOW_TABLE)

  Message: NDB Binlog: Failed to re-open shadow table %s.%s

  ER_NDB_BINLOG_CANT_REOPEN_SHADOW_TABLE was added in 8.0.2.

- Error: 4267 SQLSTATE: HY000 (ER_NDB_BINLOG_ONLINE_ALTER_RENAME_COMPLETE)

  Message: NDB Binlog: handling online alter/rename done

  ER_NDB_BINLOG_ONLINE_ALTER_RENAME_COMPLETE was added in 8.0.2.

- Error: 4268 SQLSTATE: HY000 (ER_NDB_BINLOG_SKIPPING_DROP_OF_LOCAL_TABLE)

  Message: NDB Binlog: Skipping drop of locally defined table '%s.%s' from binlog schema event '%s' from node %d.

  ER_NDB_BINLOG_SKIPPING_DROP_OF_LOCAL_TABLE was added in 8.0.2.

- Error: 4269 SQLSTATE: HY000 (ER_NDB_BINLOG_SKIPPING_RENAME_OF_LOCAL_TABLE)

  Message: NDB Binlog: Skipping renaming locally defined table '%s.%s' from binlog schema event '%s' from node %d.

  ER_NDB_BINLOG_SKIPPING_RENAME_OF_LOCAL_TABLE was added in 8.0.2.

- Error: 4270 SQLSTATE: HY000 (ER_NDB_BINLOG_SKIPPING_DROP_OF_DB_CONTAINING_LOCAL_TABLES)

  Message: NDB Binlog: Skipping drop database '%s' since it contained local tables binlog schema event '%s' from node %d.

  ER_NDB_BINLOG_SKIPPING_DROP_OF_DB_CONTAINING_LOCAL_TABLES was added in 8.0.2.

- Error: 4271 SQLSTATE: HY000 (ER_NDB_BINLOG_GOT_DIST_PRIV_EVENT_FLUSHING_PRIVILEGES)

  Message: Got dist_priv event: %s, flushing privileges

  ER_NDB_BINLOG_GOT_DIST_PRIV_EVENT_FLUSHING_PRIVILEGES was added in 8.0.2.

- Error: 4272 SQLSTATE: HY000 (ER_NDB_BINLOG_GOT_SCHEMA_EVENT)

  Message: NDB: got schema event on %s.%s(%u/%u) query: '%s' type: %s(%d) node: %u slock: %x %08x

  ER_NDB_BINLOG_GOT_SCHEMA_EVENT was added in 8.0.2.

- Error: 4273 SQLSTATE: HY000 (ER_NDB_BINLOG_SKIPPING_OLD_SCHEMA_OPERATION)

  Message: NDB schema: Skipping old schema operation(RENAME_TABLE_NEW) on %s.%s

  ER_NDB_BINLOG_SKIPPING_OLD_SCHEMA_OPERATION was added in 8.0.2.

- Error: 4274 SQLSTATE: HY000 (ER_NDB_CLUSTER_FAILURE)

  Message: NDB Schema dist: cluster failure at epoch %u/%u.

  ER_NDB_CLUSTER_FAILURE was added in 8.0.2.

- Error: 4275 SQLSTATE: HY000 (ER_NDB_TABLES_INITIALLY_READ_ONLY_ON_RECONNECT)

  Message: NDB Binlog: ndb tables initially read only on reconnect.

  ER_NDB_TABLES_INITIALLY_READ_ONLY_ON_RECONNECT was added in 8.0.2.

- Error: 4276 SQLSTATE: HY000 (ER_NDB_IGNORING_UNKNOWN_EVENT)

  Message: NDB Schema dist: unknown event %u, ignoring...

  ER_NDB_IGNORING_UNKNOWN_EVENT was added in 8.0.2.

- Error: 4277 SQLSTATE: HY000 (ER_NDB_BINLOG_OPENING_INDEX)

  Message: NDB Binlog: Opening ndb_binlog_index: %d, '%s'

  ER_NDB_BINLOG_OPENING_INDEX was added in 8.0.2.

- Error: 4278 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_LOCK_NDB_BINLOG_INDEX)

  Message: NDB Binlog: Unable to lock table ndb_binlog_index

  ER_NDB_BINLOG_CANT_LOCK_NDB_BINLOG_INDEX was added in 8.0.2.

- Error: 4279 SQLSTATE: HY000 (ER_NDB_BINLOG_INJECTING_RANDOM_WRITE_FAILURE)

  Message: NDB Binlog: Injecting random write failure

  ER_NDB_BINLOG_INJECTING_RANDOM_WRITE_FAILURE was added in 8.0.2.

- Error: 4280 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_WRITE_TO_NDB_BINLOG_INDEX)

  Message: NDB Binlog: Failed writing to ndb_binlog_index for epoch %u/%u orig_server_id %u orig_epoch %u/%u with error %d.

  ER_NDB_BINLOG_CANT_WRITE_TO_NDB_BINLOG_INDEX was added in 8.0.2.

- Error: 4281 SQLSTATE: HY000 (ER_NDB_BINLOG_WRITING_TO_NDB_BINLOG_INDEX)

  Message: NDB Binlog: Writing row (%s) to ndb_binlog_index - %s

  ER_NDB_BINLOG_WRITING_TO_NDB_BINLOG_INDEX was added in 8.0.2.

- Error: 4282 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_COMMIT_TO_NDB_BINLOG_INDEX)

  Message: NDB Binlog: Failed committing transaction to ndb_binlog_index with error %d.

  ER_NDB_BINLOG_CANT_COMMIT_TO_NDB_BINLOG_INDEX was added in 8.0.2.

- Error: 4283 SQLSTATE: HY000 (ER_NDB_BINLOG_WRITE_TO_NDB_BINLOG_INDEX_FAILED_AFTER_KILL)

Message: NDB Binlog: Failed writing to ndb_binlog_index table while retrying after kill during shutdown

ER_NDB_BINLOG_WRITE_TO_NDB_BINLOG_INDEX_FAILED_AFTER_KILL was added in 8.0.2.

- Error: 4284 SQLSTATE: HY000 (ER_NDB_BINLOG_USING_SERVER_ID_0_SLAVES_WILL_NOT)

  Message: NDB: server id set to zero - changes logged to bin log with server id zero will be logged with another server id by slave mysqlds

  ER_NDB_BINLOG_USING_SERVER_ID_0_SLAVES_WILL_NOT was added in 8.0.2.

- Error: 4285 SQLSTATE: HY000 (ER_NDB_SERVER_ID_RESERVED_OR_TOO_LARGE)

  Message: NDB: server id provided is too large to be represented in opt_server_id_bits or is reserved

  ER_NDB_SERVER_ID_RESERVED_OR_TOO_LARGE was added in 8.0.2.

- Error: 4286 SQLSTATE: HY000 (ER_NDB_BINLOG_NDB_LOG_TRANSACTION_ID_REQUIRES_V2_ROW_EVENTS)

  Message: NDB: --ndb-log-transaction-id requires v2 Binlog row events but server is using v1.

  ER_NDB_BINLOG_NDB_LOG_TRANSACTION_ID_REQUIRES_V2_ROW_EVENTS was added in 8.0.2.

- Error: 4287 SQLSTATE: HY000 (ER_NDB_BINLOG_NDB_LOG_APPLY_STATUS_FORCING_FULL_USE_WRITE)

  Message: NDB: ndb-log-apply-status forcing %s.%s to FULL USE_WRITE

  ER_NDB_BINLOG_NDB_LOG_APPLY_STATUS_FORCING_FULL_USE_WRITE was added in 8.0.2.

- Error: 4288 SQLSTATE: HY000 (ER_NDB_BINLOG_GENERIC_MESSAGE)

  Message: NDB Binlog: %s

  ER_NDB_BINLOG_GENERIC_MESSAGE was added in 8.0.2.

- Error: 4289 SQLSTATE: HY000 (ER_NDB_CONFLICT_GENERIC_MESSAGE)

  Message: %s

  ER_NDB_CONFLICT_GENERIC_MESSAGE was added in 8.0.2.

- Error: 4290 SQLSTATE: HY000 (ER_NDB_TRANS_DEPENDENCY_TRACKER_ERROR)

  Message: %s

  ER_NDB_TRANS_DEPENDENCY_TRACKER_ERROR was added in 8.0.2.

- Error: 4291 SQLSTATE: HY000 (ER_NDB_CONFLICT_FN_PARSE_ERROR)

  Message: NDB Slave: Table %s.%s : Parse error on conflict fn : %s

  ER_NDB_CONFLICT_FN_PARSE_ERROR was added in 8.0.2.

- Error: 4292 SQLSTATE: HY000 (ER_NDB_CONFLICT_FN_SETUP_ERROR)

  Message: NDB Slave: Table %s.%s : %s

  ER_NDB_CONFLICT_FN_SETUP_ERROR was added in 8.0.2.

- Error: 4293 SQLSTATE: HY000 (ER_NDB_BINLOG_FAILED_TO_GET_TABLE)

Message: NDB Binlog: Failed to get table %s from ndb: %s, %d

ER_NDB_BINLOG_FAILED_TO_GET_TABLE was added in 8.0.2.

- Error: 4294 SQLSTATE: HY000 (ER_NDB_BINLOG_NOT_LOGGING)

  Message: NDB Binlog: NOT logging %s

  ER_NDB_BINLOG_NOT_LOGGING was added in 8.0.2.

- Error: 4295 SQLSTATE: HY000 (ER_NDB_BINLOG_CREATE_TABLE_EVENT_FAILED)

  Message: NDB Binlog: FAILED CREATE (DISCOVER) TABLE Event: %s

  ER_NDB_BINLOG_CREATE_TABLE_EVENT_FAILED was added in 8.0.2.

- Error: 4296 SQLSTATE: HY000 (ER_NDB_BINLOG_CREATE_TABLE_EVENT_INFO)

  Message: NDB Binlog: CREATE (DISCOVER) TABLE Event: %s

  ER_NDB_BINLOG_CREATE_TABLE_EVENT_INFO was added in 8.0.2.

- Error: 4297 SQLSTATE: HY000 (ER_NDB_BINLOG_DISCOVER_TABLE_EVENT_INFO)

  Message: NDB Binlog: DISCOVER TABLE Event: %s

  ER_NDB_BINLOG_DISCOVER_TABLE_EVENT_INFO was added in 8.0.2.

- Error: 4298 SQLSTATE: HY000 (ER_NDB_BINLOG_BLOB_REQUIRES_PK)

  Message: NDB Binlog: logging of table %s with BLOB attribute and no PK is not supported

  ER_NDB_BINLOG_BLOB_REQUIRES_PK was added in 8.0.2.

- Error: 4299 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB)

  Message: NDB Binlog: Unable to create event in database. Event: %s Error Code: %d Message: %s

  ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB was added in 8.0.2.

- Error: 4300 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB_AND_CANT_DROP)

  Message: NDB Binlog: Unable to create event in database. Attempt to correct with drop failed. Event: %s Error Code: %d Message: %s

  ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB_AND_CANT_DROP was added in 8.0.2.

- Error: 4301 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB_DROPPED)

  Message: NDB Binlog: Unable to create event in database. Attempt to correct with drop ok, but create failed. Event: %s Error Code: %d Message: %s

  ER_NDB_BINLOG_CANT_CREATE_EVENT_IN_DB_DROPPED was added in 8.0.2.

- Error: 4302 SQLSTATE: HY000 (ER_NDB_BINLOG_DISCOVER_REUSING_OLD_EVENT_OPS)

  Message: NDB Binlog: discover reusing old ev op

  ER_NDB_BINLOG_DISCOVER_REUSING_OLD_EVENT_OPS was added in 8.0.2.

- Error: 4303 SQLSTATE: HY000 (ER_NDB_BINLOG_CREATING_NDBEVENTOPERATION_FAILED)

Message: NDB Binlog: Creating NdbEventOperation failed for %s

ER_NDB_BINLOG_CREATING_NDBEVENTOPERATION_FAILED was added in 8.0.2.

- Error: 4304 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_CREATE_BLOB)

  Message: NDB Binlog: Creating NdbEventOperation blob field %u handles failed (code=%d) for %s

  ER_NDB_BINLOG_CANT_CREATE_BLOB was added in 8.0.2.

- Error: 4305 SQLSTATE: HY000 (ER_NDB_BINLOG_NDBEVENT_EXECUTE_FAILED)

  Message: NDB Binlog: ndbevent->execute failed for %s; %d %s

  ER_NDB_BINLOG_NDBEVENT_EXECUTE_FAILED was added in 8.0.2.

- Error: 4306 SQLSTATE: HY000 (ER_NDB_CREATE_EVENT_OPS_LOGGING_INFO)

  Message: NDB Binlog: logging %s (%s,%s)

  ER_NDB_CREATE_EVENT_OPS_LOGGING_INFO was added in 8.0.2.

- Error: 4307 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_DROP_EVENT_FROM_DB)

  Message: NDB Binlog: Unable to drop event in database. Event: %s Error Code: %d Message: %s

  ER_NDB_BINLOG_CANT_DROP_EVENT_FROM_DB was added in 8.0.2.

- Error: 4308 SQLSTATE: HY000 (ER_NDB_TIMED_OUT_IN_DROP_TABLE)

  Message: NDB %s: %s timed out. Ignoring...

  ER_NDB_TIMED_OUT_IN_DROP_TABLE was added in 8.0.2.

- Error: 4309 SQLSTATE: HY000 (ER_NDB_BINLOG_UNHANDLED_ERROR_FOR_TABLE)

  Message: NDB Binlog: unhandled error %d for table %s

  ER_NDB_BINLOG_UNHANDLED_ERROR_FOR_TABLE was added in 8.0.2.

- Error: 4310 SQLSTATE: HY000 (ER_NDB_BINLOG_CLUSTER_FAILURE)

  Message: NDB Binlog: cluster failure for %s at epoch %u/%u.

  ER_NDB_BINLOG_CLUSTER_FAILURE was added in 8.0.2.

- Error: 4311 SQLSTATE: HY000 (ER_NDB_BINLOG_UNKNOWN_NON_DATA_EVENT)

  Message: NDB Binlog: unknown non data event %d for %s. Ignoring...

  ER_NDB_BINLOG_UNKNOWN_NON_DATA_EVENT was added in 8.0.2.

- Error: 4312 SQLSTATE: HY000 (ER_NDB_BINLOG_INJECTOR_DISCARDING_ROW_EVENT_METADATA)

  Message: NDB: Binlog Injector discarding row event meta data as server is using v1 row events. (%u %x)

  ER_NDB_BINLOG_INJECTOR_DISCARDING_ROW_EVENT_METADATA was added in 8.0.2.

- Error: 4313 SQLSTATE: HY000 (ER_NDB_REMAINING_OPEN_TABLES)

  Message: remove_all_event_operations: Remaining open tables:

ER_NDB_REMAINING_OPEN_TABLES was added in 8.0.2.

- Error: 4314 SQLSTATE: HY000 (ER_NDB_REMAINING_OPEN_TABLE_INFO)

  Message: %s.%s, use_count: %u

  ER_NDB_REMAINING_OPEN_TABLE_INFO was added in 8.0.2.

- Error: 4315 SQLSTATE: HY000 (ER_NDB_COULD_NOT_GET_APPLY_STATUS_SHARE)

  Message: NDB: Could not get apply status share

  ER_NDB_COULD_NOT_GET_APPLY_STATUS_SHARE was added in 8.0.2.

- Error: 4316 SQLSTATE: HY000 (ER_NDB_BINLOG_SERVER_SHUTDOWN_DURING_NDB_CLUSTER_START)

  Message: NDB Binlog: Server shutdown detected while waiting for ndbcluster to start...

  ER_NDB_BINLOG_SERVER_SHUTDOWN_DURING_NDB_CLUSTER_START was added in 8.0.2.

- Error: 4317 SQLSTATE: HY000 (ER_NDB_BINLOG_CLUSTER_RESTARTED_RESET_MASTER_SUGGESTED)

  Message: NDB Binlog: cluster has been restarted --initial or with older filesystem. ndb_latest_handled_binlog_epoch: %u/%u, while current epoch: %u/%u. RESET MASTER should be issued. Resetting ndb_latest_handled_binlog_epoch.

  ER_NDB_BINLOG_CLUSTER_RESTARTED_RESET_MASTER_SUGGESTED was added in 8.0.2.

- Error: 4318 SQLSTATE: HY000 (ER_NDB_BINLOG_CLUSTER_HAS_RECONNECTED)

  Message: NDB Binlog: cluster has reconnected. Changes to the database that occured while disconnected will not be in the binlog

  ER_NDB_BINLOG_CLUSTER_HAS_RECONNECTED was added in 8.0.2.

- Error: 4319 SQLSTATE: HY000 (ER_NDB_BINLOG_STARTING_LOG_AT_EPOCH)

  Message: NDB Binlog: starting log at epoch %u/%u

  ER_NDB_BINLOG_STARTING_LOG_AT_EPOCH was added in 8.0.2.

- Error: 4320 SQLSTATE: HY000 (ER_NDB_BINLOG_NDB_TABLES_WRITABLE)

  Message: NDB Binlog: ndb tables writable

  ER_NDB_BINLOG_NDB_TABLES_WRITABLE was added in 8.0.2.

- Error: 4321 SQLSTATE: HY000 (ER_NDB_BINLOG_SHUTDOWN_DETECTED)

  Message: NDB Binlog: Server shutdown detected...

  ER_NDB_BINLOG_SHUTDOWN_DETECTED was added in 8.0.2.

- Error: 4322 SQLSTATE: HY000 (ER_NDB_BINLOG_LOST_SCHEMA_CONNECTION_WAITING)

  Message: NDB Binlog: Just lost schema connection, hanging around

  ER_NDB_BINLOG_LOST_SCHEMA_CONNECTION_WAITING was added in 8.0.2.

- Error: 4323 SQLSTATE: HY000 (ER_NDB_BINLOG_LOST_SCHEMA_CONNECTION_CONTINUING)

Message: NDB Binlog: ...and on our way

ER_NDB_BINLOG_LOST_SCHEMA_CONNECTION_CONTINUING was added in 8.0.2.

- Error: 4324 SQLSTATE: HY000 (ER_NDB_BINLOG_ERROR_HANDLING_SCHEMA_EVENT)

  Message: NDB: error %lu (%s) on handling binlog schema event

  ER_NDB_BINLOG_ERROR_HANDLING_SCHEMA_EVENT was added in 8.0.2.

- Error: 4325 SQLSTATE: HY000 (ER_NDB_BINLOG_CANT_INJECT_APPLY_STATUS_WRITE_ROW)

  Message: NDB Binlog: Failed to inject apply status write row

  ER_NDB_BINLOG_CANT_INJECT_APPLY_STATUS_WRITE_ROW was added in 8.0.2.

- Error: 4326 SQLSTATE: HY000 (ER_NDB_BINLOG_ERROR_DURING_GCI_ROLLBACK)

  Message: NDB Binlog: Error during ROLLBACK of GCI %u/%u. Error: %d

  ER_NDB_BINLOG_ERROR_DURING_GCI_ROLLBACK was added in 8.0.2.

- Error: 4327 SQLSTATE: HY000 (ER_NDB_BINLOG_ERROR_DURING_GCI_COMMIT)

  Message: NDB Binlog: Error during COMMIT of GCI. Error: %d

  ER_NDB_BINLOG_ERROR_DURING_GCI_COMMIT was added in 8.0.2.

- Error: 4328 SQLSTATE: HY000 (ER_NDB_BINLOG_LATEST_TRX_IN_EPOCH_NOT_IN_BINLOG)

  Message: NDB Binlog: latest transaction in epoch %u/%u not in binlog as latest handled epoch is %u/%u

  ER_NDB_BINLOG_LATEST_TRX_IN_EPOCH_NOT_IN_BINLOG was added in 8.0.2.

- Error: 4329 SQLSTATE: HY000 (ER_NDB_BINLOG_RELEASING_EXTRA_SHARE_REFERENCES)

  Message: NDB Binlog: Release extra share references

  ER_NDB_BINLOG_RELEASING_EXTRA_SHARE_REFERENCES was added in 8.0.2.

- Error: 4330 SQLSTATE: HY000 (ER_NDB_BINLOG_REMAINING_OPEN_TABLES)

  Message: NDB Binlog: remaining open tables:

  ER_NDB_BINLOG_REMAINING_OPEN_TABLES was added in 8.0.2.

- Error: 4331 SQLSTATE: HY000 (ER_NDB_BINLOG_REMAINING_OPEN_TABLE_INFO)

  Message: %s.%s state: %u use_count: %u

  ER_NDB_BINLOG_REMAINING_OPEN_TABLE_INFO was added in 8.0.2.

- Error: 4332 SQLSTATE: HY000 (ER_TREE_CORRUPT_PARENT_SHOULD_POINT_AT_PARENT)

  Message: Wrong tree: Parent doesn't point at parent

  ER_TREE_CORRUPT_PARENT_SHOULD_POINT_AT_PARENT was added in 8.0.2.

- Error: 4333 SQLSTATE: HY000 (ER_TREE_CORRUPT_ROOT_SHOULD_BE_BLACK)

  Message: Wrong tree: Root should be black

  ER_TREE_CORRUPT_ROOT_SHOULD_BE_BLACK was added in 8.0.2.

- Error: 4334 SQLSTATE: HY000 (ER_TREE_CORRUPT_2_CONSECUTIVE_REDS)

  Message: Wrong tree: Found two red in a row

  ER_TREE_CORRUPT_2_CONSECUTIVE_REDS was added in 8.0.2.
- Error: 4335 SQLSTATE: HY000 (ER_TREE_CORRUPT_RIGHT_IS_LEFT)

  Message: Wrong tree: Found right == left

  ER_TREE_CORRUPT_RIGHT_IS_LEFT was added in 8.0.2.
- Error: 4336 SQLSTATE: HY000 (ER_TREE_CORRUPT_INCORRECT_BLACK_COUNT)

  Message: Wrong tree: Incorrect black-count: %d - %d

  ER_TREE_CORRUPT_INCORRECT_BLACK_COUNT was added in 8.0.2.
- Error: 4337 SQLSTATE: HY000 (ER_WRONG_COUNT_FOR_ORIGIN)

  Message: Use_count: Wrong count %lu for origin %p

  ER_WRONG_COUNT_FOR_ORIGIN was added in 8.0.2.
- Error: 4338 SQLSTATE: HY000 (ER_WRONG_COUNT_FOR_KEY)

  Message: Use_count: Wrong count for key at %p, %lu should be %lu

  ER_WRONG_COUNT_FOR_KEY was added in 8.0.2.
- Error: 4339 SQLSTATE: HY000 (ER_WRONG_COUNT_OF_ELEMENTS)

  Message: Wrong number of elements: %u (should be %u) for tree at %p

  ER_WRONG_COUNT_OF_ELEMENTS was added in 8.0.2.
- Error: 4340 SQLSTATE: HY000 (ER_RPL_ERROR_READING_SLAVE_WORKER_CONFIGURATION)

  Message: Error reading slave worker configuration

  ER_RPL_ERROR_READING_SLAVE_WORKER_CONFIGURATION was added in 8.0.2.
- Error: 4341 SQLSTATE: HY000 (ER_RPL_ERROR_WRITING_SLAVE_WORKER_CONFIGURATION)

  Message: Error writing slave worker configuration

  ER_RPL_ERROR_WRITING_SLAVE_WORKER_CONFIGURATION was added in 8.0.2.
- Error: 4342 SQLSTATE: HY000 (ER_RPL_FAILED_TO_OPEN_RELAY_LOG)

  Message: Failed to open relay log %s, error: %s

  ER_RPL_FAILED_TO_OPEN_RELAY_LOG was added in 8.0.2.
- Error: 4343 SQLSTATE: HY000 (ER_RPL_WORKER_CANT_READ_RELAY_LOG)

  Message: Error when worker read relay log events,relay log name %s, position %llu

  ER_RPL_WORKER_CANT_READ_RELAY_LOG was added in 8.0.2.
- Error: 4344 SQLSTATE: HY000 (ER_RPL_WORKER_CANT_FIND_NEXT_RELAY_LOG)

  Message: Failed to find next relay log when retrying the transaction, current relay log is %s

  ER_RPL_WORKER_CANT_FIND_NEXT_RELAY_LOG was added in 8.0.2.

- Error: 4345 SQLSTATE: HY000 (ER_RPL_MTS_SLAVE_COORDINATOR_HAS_WAITED)

  Message: Multi-threaded slave: Coordinator has waited %lu times hitting slave_pending_jobs_size_max; current event size = %zu.

  ER_RPL_MTS_SLAVE_COORDINATOR_HAS_WAITED was added in 8.0.2.

- Error: 4346 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_WRITE_DROP_FOR_TEMP_TABLES)

  Message: Failed to write the DROP statement for temporary tables to binary log

  ER_BINLOG_FAILED_TO_WRITE_DROP_FOR_TEMP_TABLES was added in 8.0.2.

- Error: 4347 SQLSTATE: HY000 (ER_BINLOG_OOM_WRITING_DELETE_WHILE_OPENING_HEAP_TABLE)

  Message: When opening HEAP table, could not allocate memory to write 'DELETE FROM `%s`.`%s`' to the binary log

  ER_BINLOG_OOM_WRITING_DELETE_WHILE_OPENING_HEAP_TABLE was added in 8.0.2.

- Error: 4348 SQLSTATE: HY000 (ER_FAILED_TO_REPAIR_TABLE)

  Message: Couldn't repair table: %s.%s

  ER_FAILED_TO_REPAIR_TABLE was added in 8.0.2.

- Error: 4349 SQLSTATE: HY000 (ER_FAILED_TO_REMOVE_TEMP_TABLE)

  Message: Could not remove temporary table: '%s', error: %d

  ER_FAILED_TO_REMOVE_TEMP_TABLE was added in 8.0.2.

- Error: 4350 SQLSTATE: HY000 (ER_SYSTEM_TABLE_NOT_TRANSACTIONAL)

  Message: System table '%.*s' is expected to be transactional.

  ER_SYSTEM_TABLE_NOT_TRANSACTIONAL was added in 8.0.2.

- Error: 4351 SQLSTATE: HY000 (ER_RPL_ERROR_WRITING_MASTER_CONFIGURATION)

  Message: Error writing master configuration.

  ER_RPL_ERROR_WRITING_MASTER_CONFIGURATION was added in 8.0.2.

- Error: 4352 SQLSTATE: HY000 (ER_RPL_ERROR_READING_MASTER_CONFIGURATION)

  Message: Error reading master configuration.

  ER_RPL_ERROR_READING_MASTER_CONFIGURATION was added in 8.0.2.

- Error: 4353 SQLSTATE: HY000 (ER_RPL_SSL_INFO_IN_MASTER_INFO_IGNORED)

  Message: SSL information in the master info file are ignored because this MySQL slave was compiled without SSL support.

  ER_RPL_SSL_INFO_IN_MASTER_INFO_IGNORED was added in 8.0.2.

- Error: 4354 SQLSTATE: HY000 (ER_PLUGIN_FAILED_DEINITIALIZATION)

  Message: Plugin '%s' of type %s failed deinitialization

  ER_PLUGIN_FAILED_DEINITIALIZATION was added in 8.0.2.

- Error: 4355 SQLSTATE: HY000 (ER_PLUGIN_HAS_NONZERO_REFCOUNT_AFTER_DEINITIALIZATION)

  Message: Plugin '%s' has ref_count=%d after deinitialization.

  ER_PLUGIN_HAS_NONZERO_REFCOUNT_AFTER_DEINITIALIZATION was added in 8.0.2.

- Error: 4356 SQLSTATE: HY000 (ER_PLUGIN_SHUTTING_DOWN_PLUGIN)

  Message: Shutting down plugin '%s'

  ER_PLUGIN_SHUTTING_DOWN_PLUGIN was added in 8.0.2.

- Error: 4357 SQLSTATE: HY000 (ER_PLUGIN_REGISTRATION_FAILED)

  Message: Plugin '%s' registration as a %s failed.

  ER_PLUGIN_REGISTRATION_FAILED was added in 8.0.2.

- Error: 4358 SQLSTATE: HY000 (ER_PLUGIN_CANT_OPEN_PLUGIN_TABLE)

  Message: Can't open the mysql.plugin table. Please run mysql_upgrade to create it.

  ER_PLUGIN_CANT_OPEN_PLUGIN_TABLE was added in 8.0.2.

- Error: 4359 SQLSTATE: HY000 (ER_PLUGIN_CANT_LOAD)

  Message: Couldn't load plugin named '%s' with soname '%s'.

  ER_PLUGIN_CANT_LOAD was added in 8.0.2.

- Error: 4360 SQLSTATE: HY000 (ER_PLUGIN_LOAD_PARAMETER_TOO_LONG)

  Message: plugin-load parameter too long

  ER_PLUGIN_LOAD_PARAMETER_TOO_LONG was added in 8.0.2.

- Error: 4361 SQLSTATE: HY000 (ER_PLUGIN_FORCING_SHUTDOWN)

  Message: Plugin '%s' will be forced to shutdown

  ER_PLUGIN_FORCING_SHUTDOWN was added in 8.0.2.

- Error: 4362 SQLSTATE: HY000 (ER_PLUGIN_HAS_NONZERO_REFCOUNT_AFTER_SHUTDOWN)

  Message: Plugin '%s' has ref_count=%d after shutdown.

  ER_PLUGIN_HAS_NONZERO_REFCOUNT_AFTER_SHUTDOWN was added in 8.0.2.

- Error: 4363 SQLSTATE: HY000 (ER_PLUGIN_UNKNOWN_VARIABLE_TYPE)

  Message: Unknown variable type code 0x%x in plugin '%s'.

  ER_PLUGIN_UNKNOWN_VARIABLE_TYPE was added in 8.0.2.

- Error: 4364 SQLSTATE: HY000 (ER_PLUGIN_VARIABLE_SET_READ_ONLY)

  Message: Server variable %s of plugin %s was forced to be read-only: string variable without update_func and PLUGIN_VAR_MEMALLOC flag

  ER_PLUGIN_VARIABLE_SET_READ_ONLY was added in 8.0.2.

- Error: 4365 SQLSTATE: HY000 (ER_PLUGIN_VARIABLE_MISSING_NAME)

Message: Missing variable name in plugin '%s'.

ER_PLUGIN_VARIABLE_MISSING_NAME was added in 8.0.2.

- Error: 4366 SQLSTATE: HY000 (ER_PLUGIN_VARIABLE_NOT_ALLOCATED_THREAD_LOCAL)

  Message: Thread local variable '%s' not allocated in plugin '%s'.

  ER_PLUGIN_VARIABLE_NOT_ALLOCATED_THREAD_LOCAL was added in 8.0.2.

- Error: 4367 SQLSTATE: HY000 (ER_PLUGIN_OOM)

  Message: Out of memory for plugin '%s'.

  ER_PLUGIN_OOM was added in 8.0.2.

- Error: 4368 SQLSTATE: HY000 (ER_PLUGIN_BAD_OPTIONS)

  Message: Bad options for plugin '%s'.

  ER_PLUGIN_BAD_OPTIONS was added in 8.0.2.

- Error: 4369 SQLSTATE: HY000 (ER_PLUGIN_PARSING_OPTIONS_FAILED)

  Message: Parsing options for plugin '%s' failed.

  ER_PLUGIN_PARSING_OPTIONS_FAILED was added in 8.0.2.

- Error: 4370 SQLSTATE: HY000 (ER_PLUGIN_DISABLED)

  Message: Plugin '%s' is disabled.

  ER_PLUGIN_DISABLED was added in 8.0.2.

- Error: 4371 SQLSTATE: HY000 (ER_PLUGIN_HAS_CONFLICTING_SYSTEM_VARIABLES)

  Message: Plugin '%s' has conflicting system variables

  ER_PLUGIN_HAS_CONFLICTING_SYSTEM_VARIABLES was added in 8.0.2.

- Error: 4372 SQLSTATE: HY000 (ER_PLUGIN_CANT_SET_PERSISTENT_OPTIONS)

  Message: Setting persistent options for plugin '%s' failed.

  ER_PLUGIN_CANT_SET_PERSISTENT_OPTIONS was added in 8.0.2.

- Error: 4373 SQLSTATE: HY000 (ER_MY_NET_WRITE_FAILED_FALLING_BACK_ON_STDERR)

  Message: Failed on my_net_write, writing to stderr instead: %s

  ER_MY_NET_WRITE_FAILED_FALLING_BACK_ON_STDERR was added in 8.0.2.

- Error: 4374 SQLSTATE: HY000 (ER_RETRYING_REPAIR_WITHOUT_QUICK)

  Message: Retrying repair of: '%s' without quick

  ER_RETRYING_REPAIR_WITHOUT_QUICK was added in 8.0.2.

- Error: 4375 SQLSTATE: HY000 (ER_RETRYING_REPAIR_WITH_KEYCACHE)

  Message: Retrying repair of: '%s' with keycache

  ER_RETRYING_REPAIR_WITH_KEYCACHE was added in 8.0.2.

- Error: 4376 SQLSTATE: HY000 (ER_FOUND_ROWS_WHILE_REPAIRING)

  Message: Found %s of %s rows when repairing '%s'

  ER_FOUND_ROWS_WHILE_REPAIRING was added in 8.0.2.
- Error: 4377 SQLSTATE: HY000 (ER_ERROR_DURING_OPTIMIZE_TABLE)

  Message: Warning: Optimize table got errno %d on %s.%s, retrying

  ER_ERROR_DURING_OPTIMIZE_TABLE was added in 8.0.2.
- Error: 4378 SQLSTATE: HY000 (ER_ERROR_ENABLING_KEYS)

  Message: Warning: Enabling keys got errno %d on %s.%s, retrying

  ER_ERROR_ENABLING_KEYS was added in 8.0.2.
- Error: 4379 SQLSTATE: HY000 (ER_CHECKING_TABLE)

  Message: Checking table: '%s'

  ER_CHECKING_TABLE was added in 8.0.2.
- Error: 4380 SQLSTATE: HY000 (ER_RECOVERING_TABLE)

  Message: Recovering table: '%s'

  ER_RECOVERING_TABLE was added in 8.0.2.
- Error: 4381 SQLSTATE: HY000 (ER_CANT_CREATE_TABLE_SHARE_FROM_FRM)

  Message: Error in creating TABLE_SHARE from %s.frm file.

  ER_CANT_CREATE_TABLE_SHARE_FROM_FRM was added in 8.0.2.
- Error: 4382 SQLSTATE: HY000 (ER_CANT_LOCK_TABLE)

  Message: Unable to acquire lock on %s.%s

  ER_CANT_LOCK_TABLE was added in 8.0.2.
- Error: 4383 SQLSTATE: HY000 (ER_CANT_ALLOC_TABLE_OBJECT)

  Message: Error in allocation memory for TABLE object.

  ER_CANT_ALLOC_TABLE_OBJECT was added in 8.0.2.
- Error: 4384 SQLSTATE: HY000 (ER_CANT_CREATE_HANDLER_OBJECT_FOR_TABLE)

  Message: Error in creating handler object for table %s.%s

  ER_CANT_CREATE_HANDLER_OBJECT_FOR_TABLE was added in 8.0.2.
- Error: 4385 SQLSTATE: HY000 (ER_CANT_SET_HANDLER_REFERENCE_FOR_TABLE)

  Message: Error in setting handler reference for table %s.%s

  ER_CANT_SET_HANDLER_REFERENCE_FOR_TABLE was added in 8.0.2.
- Error: 4386 SQLSTATE: HY000 (ER_CANT_LOCK_TABLESPACE)

  Message: Unable to acquire lock on tablespace name %s

  ER_CANT_LOCK_TABLESPACE was added in 8.0.2.

- Error: 4387 SQLSTATE: HY000 (ER_CANT_UPGRADE_GENERATED_COLUMNS_TO_DD)

  Message: Error in processing generated columns

  ER_CANT_UPGRADE_GENERATED_COLUMNS_TO_DD was added in 8.0.2.
- Error: 4388 SQLSTATE: HY000 (ER_DD_ERROR_CREATING_ENTRY)

  Message: Error in Creating DD entry for %s.%s

  ER_DD_ERROR_CREATING_ENTRY was added in 8.0.2.
- Error: 4389 SQLSTATE: HY000 (ER_DD_CANT_FETCH_TABLE_DATA)

  Message: Error in fetching %s.%s table data from dictionary

  ER_DD_CANT_FETCH_TABLE_DATA was added in 8.0.2.
- Error: 4390 SQLSTATE: HY000 (ER_DD_CANT_FIX_SE_DATA)

  Message: Error in fixing SE data for %s.%s

  ER_DD_CANT_FIX_SE_DATA was added in 8.0.2.
- Error: 4391 SQLSTATE: HY000 (ER_DD_CANT_CREATE_SP)

  Message: Error in creating stored program '%s.%s'

  ER_DD_CANT_CREATE_SP was added in 8.0.2.
- Error: 4392 SQLSTATE: HY000 (ER_CANT_OPEN_DB_OPT_USING_DEFAULT_CHARSET)

  Message: Unable to open db.opt file %s. Using default Character set.

  ER_CANT_OPEN_DB_OPT_USING_DEFAULT_CHARSET was added in 8.0.2.
- Error: 4393 SQLSTATE: HY000 (ER_CANT_CREATE_CACHE_FOR_DB_OPT)

  Message: Unable to intialize IO cache to open db.opt file %s.

  ER_CANT_CREATE_CACHE_FOR_DB_OPT was added in 8.0.2.
- Error: 4394 SQLSTATE: HY000 (ER_CANT_IDENTIFY_CHARSET_USING_DEFAULT)

  Message: Unable to identify the charset in %s. Using default character set.

  ER_CANT_IDENTIFY_CHARSET_USING_DEFAULT was added in 8.0.2.
- Error: 4395 SQLSTATE: HY000 (ER_DB_OPT_NOT_FOUND_USING_DEFAULT_CHARSET)

  Message: db.opt file not found for %s database. Using default Character set.

  ER_DB_OPT_NOT_FOUND_USING_DEFAULT_CHARSET was added in 8.0.2.
- Error: 4396 SQLSTATE: HY000 (ER_EVENT_CANT_GET_TIMEZONE_FROM_FIELD)

  Message: Event '%s'.'%s': invalid value in column mysql.event.time_zone.

  ER_EVENT_CANT_GET_TIMEZONE_FROM_FIELD was added in 8.0.2.
- Error: 4397 SQLSTATE: HY000 (ER_EVENT_CANT_FIND_TIMEZONE)

  Message: Event '%s'.'%s': has invalid time zone value

  ER_EVENT_CANT_FIND_TIMEZONE was added in 8.0.2.

- Error: 4398 SQLSTATE: HY000 (ER_EVENT_CANT_GET_CHARSET)

  Message: Event '%s'.'%s': invalid value in column mysql.event.character_set_client.

  ER_EVENT_CANT_GET_CHARSET was added in 8.0.2.

- Error: 4399 SQLSTATE: HY000 (ER_EVENT_CANT_GET_COLLATION)

  Message: Event '%s'.'%s': invalid value in column mysql.event.collation_connection.

  ER_EVENT_CANT_GET_COLLATION was added in 8.0.2.

- Error: 4400 SQLSTATE: HY000 (ER_EVENT_CANT_OPEN_TABLE_MYSQL_EVENT)

  Message: Failed to open mysql.event Table.

  ER_EVENT_CANT_OPEN_TABLE_MYSQL_EVENT was added in 8.0.2.

- Error: 4401 SQLSTATE: HY000 (ER_CANT_PARSE_STORED_ROUTINE_BODY)

  Message: Parsing '%s.%s' routine body failed. Creating routine without parsing routine body

  ER_CANT_PARSE_STORED_ROUTINE_BODY was added in 8.0.2.

- Error: 4402 SQLSTATE: HY000 (ER_CANT_OPEN_TABLE_MYSQL_PROC)

  Message: Failed to open mysql.proc Table.

  ER_CANT_OPEN_TABLE_MYSQL_PROC was added in 8.0.2.

- Error: 4403 SQLSTATE: HY000 (ER_CANT_READ_TABLE_MYSQL_PROC)

  Message: Failed to read mysql.proc table.

  ER_CANT_READ_TABLE_MYSQL_PROC was added in 8.0.2.

- Error: 4404 SQLSTATE: HY000 (ER_FILE_EXISTS_DURING_UPGRADE)

  Message: Found %s file in mysql schema. DD will create .ibd file with same name. Please rename table and start upgrade process again.

  ER_FILE_EXISTS_DURING_UPGRADE was added in 8.0.2.

- Error: 4405 SQLSTATE: HY000 (ER_CANT_OPEN_DATADIR_AFTER_UPGRADE_FAILURE)

  Message: Unable to open the data directory %s during clean up after upgrade failed

  ER_CANT_OPEN_DATADIR_AFTER_UPGRADE_FAILURE was added in 8.0.2.

- Error: 4406 SQLSTATE: HY000 (ER_CANT_SET_PATH_FOR)

  Message: Failed to set path %s

  ER_CANT_SET_PATH_FOR was added in 8.0.2.

- Error: 4407 SQLSTATE: HY000 (ER_CANT_OPEN_DIR)

  Message: Failed to open dir %s

  ER_CANT_OPEN_DIR was added in 8.0.2.

- Error: 4408 SQLSTATE: HY000 (ER_NDB_EMPTY_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

Message: NDB: Found empty nodeid specified in --ndb-cluster-connection-pool-nodeids='%s'.

ER_NDB_EMPTY_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4409 SQLSTATE: HY000 (ER_NDB_CANT_PARSE_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

  Message: NDB: Could not parse '%s' in --ndb-cluster-connection-pool-nodeids='%s'.

  ER_NDB_CANT_PARSE_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4410 SQLSTATE: HY000 (ER_NDB_INVALID_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

  Message: NDB: Invalid nodeid %d in --ndb-cluster-connection-pool-nodeids='%s'.

  ER_NDB_INVALID_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4411 SQLSTATE: HY000 (ER_NDB_DUPLICATE_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

  Message: NDB: Found duplicate nodeid %d in --ndb-cluster-connection-pool-nodeids='%s'.

  ER_NDB_DUPLICATE_NODEID_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4412 SQLSTATE: HY000 (ER_NDB_POOL_SIZE_MUST_MATCH_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

  Message: NDB: The size of the cluster connection pool must be equal to the number of nodeids in --ndb-cluster-connection-pool-nodeids='%s'.

  ER_NDB_POOL_SIZE_MUST_MATCH_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4413 SQLSTATE: HY000 (ER_NDB_NODEID_NOT_FIRST_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS)

  Message: NDB: The nodeid specified by --ndb-nodeid must be equal to the first nodeid in --ndb-cluster-connection-pool-nodeids='%s'.

  ER_NDB_NODEID_NOT_FIRST_IN_NDB_CLUSTER_CONNECTION_POOL_NODEIDS was added in 8.0.2.

- Error: 4414 SQLSTATE: HY000 (ER_NDB_USING_NODEID)

  Message: NDB: using nodeid %u

  ER_NDB_USING_NODEID was added in 8.0.2.

- Error: 4415 SQLSTATE: HY000 (ER_NDB_CANT_ALLOC_GLOBAL_NDB_CLUSTER_CONNECTION)

  Message: NDB: failed to allocate global ndb cluster connection

  ER_NDB_CANT_ALLOC_GLOBAL_NDB_CLUSTER_CONNECTION was added in 8.0.2.

- Error: 4416 SQLSTATE: HY000 (ER_NDB_CANT_ALLOC_GLOBAL_NDB_OBJECT)

  Message: NDB: failed to allocate global ndb object

  ER_NDB_CANT_ALLOC_GLOBAL_NDB_OBJECT was added in 8.0.2.

- Error: 4417 SQLSTATE: HY000 (ER_NDB_USING_NODEID_LIST)

  Message: NDB[%u]: using nodeid %u

  ER_NDB_USING_NODEID_LIST was added in 8.0.2.

- Error: 4418 SQLSTATE: HY000 (ER_NDB_CANT_ALLOC_NDB_CLUSTER_CONNECTION)

  Message: NDB[%u]: failed to allocate cluster connect object

  ER_NDB_CANT_ALLOC_NDB_CLUSTER_CONNECTION was added in 8.0.2.

- Error: 4419 SQLSTATE: HY000 (ER_NDB_STARTING_CONNECT_THREAD)

  Message: NDB[%u]: starting connect thread

  ER_NDB_STARTING_CONNECT_THREAD was added in 8.0.2.

- Error: 4420 SQLSTATE: HY000 (ER_NDB_NODE_INFO)

  Message: NDB[%u]: NodeID: %d, %s

  ER_NDB_NODE_INFO was added in 8.0.2.

- Error: 4421 SQLSTATE: HY000 (ER_NDB_CANT_START_CONNECT_THREAD)

  Message: NDB[%u]: failed to start connect thread

  ER_NDB_CANT_START_CONNECT_THREAD was added in 8.0.2.

- Error: 4422 SQLSTATE: HY000 (ER_NDB_GENERIC_ERROR)

  Message: NDB: error (%u) %s

  ER_NDB_GENERIC_ERROR was added in 8.0.2.

- Error: 4423 SQLSTATE: HY000 (ER_NDB_CPU_MASK_TOO_SHORT)

  Message: Ignored receive thread CPU mask, mask too short, %u CPUs needed in mask, only %u CPUs provided

  ER_NDB_CPU_MASK_TOO_SHORT was added in 8.0.2.

- Error: 4424 SQLSTATE: HY000 (ER_EVENT_ERROR_CREATING_QUERY_TO_WRITE_TO_BINLOG)

  Message: Event Error: An error occurred while creating query string, before writing it into binary log.

  ER_EVENT_ERROR_CREATING_QUERY_TO_WRITE_TO_BINLOG was added in 8.0.2.

- Error: 4425 SQLSTATE: HY000 (ER_EVENT_SCHEDULER_ERROR_LOADING_FROM_DB)

  Message: Event Scheduler: Error while loading from disk.

  ER_EVENT_SCHEDULER_ERROR_LOADING_FROM_DB was added in 8.0.2.

- Error: 4426 SQLSTATE: HY000 (ER_EVENT_SCHEDULER_ERROR_GETTING_EVENT_OBJECT)

  Message: Event Scheduler: Error getting event object.

  ER_EVENT_SCHEDULER_ERROR_GETTING_EVENT_OBJECT was added in 8.0.2.

- Error: 4427 SQLSTATE: HY000 (ER_EVENT_SCHEDULER_GOT_BAD_DATA_FROM_TABLE)

Message: Event Scheduler: Error while loading events from mysql.events.The table probably contains bad data or is corrupted

`ER_EVENT_SCHEDULER_GOT_BAD_DATA_FROM_TABLE` was added in 8.0.2.

- Error: `4428` SQLSTATE: `HY000` (`ER_EVENT_CANT_GET_LOCK_FOR_DROPPING_EVENT`)

  Message: Unable to obtain lock for dropping event %s from schema %s

  `ER_EVENT_CANT_GET_LOCK_FOR_DROPPING_EVENT` was added in 8.0.2.

- Error: `4429` SQLSTATE: `HY000` (`ER_EVENT_UNABLE_TO_DROP_EVENT`)

  Message: Unable to drop event %s from schema %s

  `ER_EVENT_UNABLE_TO_DROP_EVENT` was added in 8.0.2.

- Error: `4430` SQLSTATE: `HY000` (`ER_BINLOG_ATTACHING_THREAD_MEMORY_FINALLY_AVAILABLE`)

  Message: Server overcomes the temporary 'out of memory' in '%d' tries while attaching to session thread during the group commit phase.

  `ER_BINLOG_ATTACHING_THREAD_MEMORY_FINALLY_AVAILABLE` was added in 8.0.2.

- Error: `4431` SQLSTATE: `HY000` (`ER_BINLOG_CANT_RESIZE_CACHE`)

  Message: Unable to resize binlog IOCACHE auxiliary file

  `ER_BINLOG_CANT_RESIZE_CACHE` was added in 8.0.2.

- Error: `4432` SQLSTATE: `HY000` (`ER_BINLOG_FILE_BEING_READ_NOT_PURGED`)

  Message: file %s was not purged because it was being read by thread number %u

  `ER_BINLOG_FILE_BEING_READ_NOT_PURGED` was added in 8.0.2.

- Error: `4433` SQLSTATE: `HY000` (`ER_BINLOG_IO_ERROR_READING_HEADER`)

  Message: I/O error reading the header from the binary log, errno=%d, io cache code=%d

  `ER_BINLOG_IO_ERROR_READING_HEADER` was added in 8.0.2.

- Error: `4434` SQLSTATE: `HY000` (`ER_BINLOG_CANT_OPEN_LOG`)

  Message: Failed to open log (file '%s', errno %d)

  `ER_BINLOG_CANT_OPEN_LOG` was added in 8.0.2.

- Error: `4435` SQLSTATE: `HY000` (`ER_BINLOG_CANT_CREATE_CACHE_FOR_LOG`)

  Message: Failed to create a cache on log (file '%s')

  `ER_BINLOG_CANT_CREATE_CACHE_FOR_LOG` was added in 8.0.2.

- Error: `4436` SQLSTATE: `HY000` (`ER_BINLOG_FILE_EXTENSION_NUMBER_EXHAUSTED`)

  Message: Log filename extension number exhausted: %06lu. Please fix this by archiving old logs and updating the index files.

  `ER_BINLOG_FILE_EXTENSION_NUMBER_EXHAUSTED` was added in 8.0.2.

- Error: `4437` SQLSTATE: `HY000` (`ER_BINLOG_FILE_NAME_TOO_LONG`)

Message: Log filename too large: %s%s (%zu). Please fix this by archiving old logs and updating the index files.

ER_BINLOG_FILE_NAME_TOO_LONG was added in 8.0.2.

- Error: 4438 SQLSTATE: HY000 (ER_BINLOG_FILE_EXTENSION_NUMBER_RUNNING_LOW)

  Message: Next log extension: %lu. Remaining log filename extensions: %lu. Please consider archiving some logs.

  ER_BINLOG_FILE_EXTENSION_NUMBER_RUNNING_LOW was added in 8.0.2.

- Error: 4439 SQLSTATE: HY000 (ER_BINLOG_CANT_OPEN_FOR_LOGGING)

  Message: Could not open %s for logging (error %d). Turning logging off for the whole duration of the MySQL server process. To turn it on again: fix the cause, shutdown the MySQL server and restart it.

  ER_BINLOG_CANT_OPEN_FOR_LOGGING was added in 8.0.2.

- Error: 4440 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_SYNC_INDEX_FILE)

  Message: MYSQL_BIN_LOG::open_index_file failed to sync the index file.

  ER_BINLOG_FAILED_TO_SYNC_INDEX_FILE was added in 8.0.2.

- Error: 4441 SQLSTATE: HY000 (ER_BINLOG_ERROR_READING_GTIDS_FROM_RELAY_LOG)

  Message: Error reading GTIDs from relaylog: %d

  ER_BINLOG_ERROR_READING_GTIDS_FROM_RELAY_LOG was added in 8.0.2.

- Error: 4442 SQLSTATE: HY000 (ER_BINLOG_EVENTS_READ_FROM_RELAY_LOG_INFO)

  Message: %lu events read in relaylog file '%s' for updating Retrieved_Gtid_Set and/or IO thread transaction parser state.

  ER_BINLOG_EVENTS_READ_FROM_RELAY_LOG_INFO was added in 8.0.2.

- Error: 4443 SQLSTATE: HY000 (ER_BINLOG_ERROR_READING_GTIDS_FROM_BINARY_LOG)

  Message: Error reading GTIDs from binary log: %d

  ER_BINLOG_ERROR_READING_GTIDS_FROM_BINARY_LOG was added in 8.0.2.

- Error: 4444 SQLSTATE: HY000 (ER_BINLOG_EVENTS_READ_FROM_BINLOG_INFO)

  Message: Read %lu events from binary log file '%s' to determine the GTIDs purged from binary logs.

  ER_BINLOG_EVENTS_READ_FROM_BINLOG_INFO was added in 8.0.2.

- Error: 4445 SQLSTATE: HY000 (ER_BINLOG_CANT_GENERATE_NEW_FILE_NAME)

  Message: MYSQL_BIN_LOG::open failed to generate new file name.

  ER_BINLOG_CANT_GENERATE_NEW_FILE_NAME was added in 8.0.2.

- Error: 4446 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_SYNC_INDEX_FILE_IN_OPEN)

  Message: MYSQL_BIN_LOG::open failed to sync the index file.

  ER_BINLOG_FAILED_TO_SYNC_INDEX_FILE_IN_OPEN was added in 8.0.2.

- Error: 4447 SQLSTATE: HY000 (ER_BINLOG_CANT_USE_FOR_LOGGING)

Message: Could not use %s for logging (error %d). Turning logging off for the whole duration of the MySQL server process. To turn it on again: fix the cause, shutdown the MySQL server and restart it.

`ER_BINLOG_CANT_USE_FOR_LOGGING` was added in 8.0.2.

- Error: `4448` SQLSTATE: `HY000` (`ER_BINLOG_FAILED_TO_CLOSE_INDEX_FILE_WHILE_REBUILDING`)

  Message: While rebuilding index file %s: Failed to close the index file.

  `ER_BINLOG_FAILED_TO_CLOSE_INDEX_FILE_WHILE_REBUILDING` was added in 8.0.2.

- Error: `4449` SQLSTATE: `HY000` (`ER_BINLOG_FAILED_TO_DELETE_INDEX_FILE_WHILE_REBUILDING`)

  Message: While rebuilding index file %s: Failed to delete the existing index file. It could be that file is being used by some other process.

  `ER_BINLOG_FAILED_TO_DELETE_INDEX_FILE_WHILE_REBUILDING` was added in 8.0.2.

- Error: `4450` SQLSTATE: `HY000` (`ER_BINLOG_FAILED_TO_RENAME_INDEX_FILE_WHILE_REBUILDING`)

  Message: While rebuilding index file %s: Failed to rename the new index file to the existing index file.

  `ER_BINLOG_FAILED_TO_RENAME_INDEX_FILE_WHILE_REBUILDING` was added in 8.0.2.

- Error: `4451` SQLSTATE: `HY000` (`ER_BINLOG_FAILED_TO_OPEN_INDEX_FILE_AFTER_REBUILDING`)

  Message: After rebuilding the index file %s: Failed to open the index file.

  `ER_BINLOG_FAILED_TO_OPEN_INDEX_FILE_AFTER_REBUILDING` was added in 8.0.2.

- Error: `4452` SQLSTATE: `HY000` (`ER_BINLOG_CANT_APPEND_LOG_TO_TMP_INDEX`)

  Message: MYSQL_BIN_LOG::add_log_to_index failed to append log file name: %s, to crash safe index file.

  `ER_BINLOG_CANT_APPEND_LOG_TO_TMP_INDEX` was added in 8.0.2.

- Error: `4453` SQLSTATE: `HY000` (`ER_BINLOG_CANT_LOCATE_OLD_BINLOG_OR_RELAY_LOG_FILES`)

  Message: Failed to locate old binlog or relay log files

  `ER_BINLOG_CANT_LOCATE_OLD_BINLOG_OR_RELAY_LOG_FILES` was added in 8.0.2.

- Error: `4454` SQLSTATE: `HY000` (`ER_BINLOG_CANT_DELETE_FILE`)

  Message: Failed to delete file '%s'

  `ER_BINLOG_CANT_DELETE_FILE` was added in 8.0.2.

- Error: `4455` SQLSTATE: `HY000` (`ER_BINLOG_CANT_SET_TMP_INDEX_NAME`)

  Message: MYSQL_BIN_LOG::set_crash_safe_index_file_name failed to set file name.

  `ER_BINLOG_CANT_SET_TMP_INDEX_NAME` was added in 8.0.2.

- Error: `4456` SQLSTATE: `HY000` (`ER_BINLOG_FAILED_TO_OPEN_TEMPORARY_INDEX_FILE`)

  Message: MYSQL_BIN_LOG::open_crash_safe_index_file failed to open temporary index file.

ER_BINLOG_FAILED_TO_OPEN_TEMPORARY_INDEX_FILE was added in 8.0.2.

- Error: 4457 SQLSTATE: HY000 (ER_BINLOG_ERROR_GETTING_NEXT_LOG_FROM_INDEX)

  Message: next log error: %d offset: %s log: %s included: %d

  ER_BINLOG_ERROR_GETTING_NEXT_LOG_FROM_INDEX was added in 8.0.2.

- Error: 4458 SQLSTATE: HY000 (ER_BINLOG_CANT_OPEN_TMP_INDEX)

  Message: %s failed to open the crash safe index file.

  ER_BINLOG_CANT_OPEN_TMP_INDEX was added in 8.0.2.

- Error: 4459 SQLSTATE: HY000 (ER_BINLOG_CANT_COPY_INDEX_TO_TMP)

  Message: %s failed to copy index file to crash safe index file.

  ER_BINLOG_CANT_COPY_INDEX_TO_TMP was added in 8.0.2.

- Error: 4460 SQLSTATE: HY000 (ER_BINLOG_CANT_CLOSE_TMP_INDEX)

  Message: %s failed to close the crash safe index file.

  ER_BINLOG_CANT_CLOSE_TMP_INDEX was added in 8.0.2.

- Error: 4461 SQLSTATE: HY000 (ER_BINLOG_CANT_MOVE_TMP_TO_INDEX)

  Message: %s failed to move crash safe index file to index file.

  ER_BINLOG_CANT_MOVE_TMP_TO_INDEX was added in 8.0.2.

- Error: 4462 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_CALLED_WITH_FILE_NOT_IN_INDEX)

  Message: MYSQL_BIN_LOG::purge_logs was called with file %s not listed in the index.

  ER_BINLOG_PURGE_LOGS_CALLED_WITH_FILE_NOT_IN_INDEX was added in 8.0.2.

- Error: 4463 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_CANT_SYNC_INDEX_FILE)

  Message: MYSQL_BIN_LOG::purge_logs failed to sync the index file.

  ER_BINLOG_PURGE_LOGS_CANT_SYNC_INDEX_FILE was added in 8.0.2.

- Error: 4464 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_CANT_COPY_TO_REGISTER_FILE)

  Message: MYSQL_BIN_LOG::purge_logs failed to copy %s to register file.

  ER_BINLOG_PURGE_LOGS_CANT_COPY_TO_REGISTER_FILE was added in 8.0.2.

- Error: 4465 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_CANT_FLUSH_REGISTER_FILE)

  Message: MYSQL_BIN_LOG::purge_logs failed to flush register file.

  ER_BINLOG_PURGE_LOGS_CANT_FLUSH_REGISTER_FILE was added in 8.0.2.

- Error: 4466 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_CANT_UPDATE_INDEX_FILE)

  Message: MYSQL_BIN_LOG::purge_logs failed to update the index file

  ER_BINLOG_PURGE_LOGS_CANT_UPDATE_INDEX_FILE was added in 8.0.2.

- Error: 4467 SQLSTATE: HY000 (ER_BINLOG_PURGE_LOGS_FAILED_TO_PURGE_LOG)

Message: MYSQL_BIN_LOG::purge_logs failed to process registered files that would be purged.

ER_BINLOG_PURGE_LOGS_FAILED_TO_PURGE_LOG was added in 8.0.2.

- Error: 4468 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_SET_PURGE_INDEX_FILE_NAME)

  Message: MYSQL_BIN_LOG::set_purge_index_file_name failed to set file name.

  ER_BINLOG_FAILED_TO_SET_PURGE_INDEX_FILE_NAME was added in 8.0.2.

- Error: 4469 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_OPEN_REGISTER_FILE)

  Message: MYSQL_BIN_LOG::open_purge_index_file failed to open register file.

  ER_BINLOG_FAILED_TO_OPEN_REGISTER_FILE was added in 8.0.2.

- Error: 4470 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_REINIT_REGISTER_FILE)

  Message: MYSQL_BIN_LOG::purge_index_entry failed to reinit register file for read

  ER_BINLOG_FAILED_TO_REINIT_REGISTER_FILE was added in 8.0.2.

- Error: 4471 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_READ_REGISTER_FILE)

  Message: MYSQL_BIN_LOG::purge_index_entry error %d reading from register file.

  ER_BINLOG_FAILED_TO_READ_REGISTER_FILE was added in 8.0.2.

- Error: 4472 SQLSTATE: HY000 (ER_CANT_STAT_FILE)

  Message: Failed to execute mysql_file_stat on file '%s'

  ER_CANT_STAT_FILE was added in 8.0.2.

- Error: 4473 SQLSTATE: HY000 (ER_BINLOG_CANT_DELETE_LOG_FILE_DOES_INDEX_MATCH_FILES)

  Message: Failed to delete log file '%s'; consider examining correspondence of your binlog index file to the actual binlog files

  ER_BINLOG_CANT_DELETE_LOG_FILE_DOES_INDEX_MATCH_FILES was added in 8.0.2.

- Error: 4474 SQLSTATE: HY000 (ER_BINLOG_CANT_DELETE_FILE_AND_READ_BINLOG_INDEX)

  Message: Failed to delete file '%s' and read the binlog index file

  ER_BINLOG_CANT_DELETE_FILE_AND_READ_BINLOG_INDEX was added in 8.0.2.

- Error: 4475 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_DELETE_LOG_FILE)

  Message: Failed to delete log file '%s'

  ER_BINLOG_FAILED_TO_DELETE_LOG_FILE was added in 8.0.2.

- Error: 4476 SQLSTATE: HY000 (ER_BINLOG_LOGGING_INCIDENT_TO_STOP_SLAVES)

  Message: %s An incident event has been written to the binary log which will stop the slaves.

  ER_BINLOG_LOGGING_INCIDENT_TO_STOP_SLAVES was added in 8.0.2.

- Error: 4477 SQLSTATE: HY000 (ER_BINLOG_CANT_FIND_LOG_IN_INDEX)

  Message: find_log_pos() failed (error: %d)

ER_BINLOG_CANT_FIND_LOG_IN_INDEX was added in 8.0.2.

- Error: 4478 SQLSTATE: HY000 (ER_BINLOG_RECOVERING_AFTER_CRASH_USING)

  Message: Recovering after a crash using %s

  ER_BINLOG_RECOVERING_AFTER_CRASH_USING was added in 8.0.2.

- Error: 4479 SQLSTATE: HY000 (ER_BINLOG_CANT_OPEN_CRASHED_BINLOG)

  Message: Failed to open the crashed binlog file when master server is recovering it.

  ER_BINLOG_CANT_OPEN_CRASHED_BINLOG was added in 8.0.2.

- Error: 4480 SQLSTATE: HY000 (ER_BINLOG_CANT_TRIM_CRASHED_BINLOG)

  Message: Failed to trim the crashed binlog file when master server is recovering it.

  ER_BINLOG_CANT_TRIM_CRASHED_BINLOG was added in 8.0.2.

- Error: 4481 SQLSTATE: HY000 (ER_BINLOG_CRASHED_BINLOG_TRIMMED)

  Message: Crashed binlog file %s size is %llu, but recovered up to %llu. Binlog trimmed to %llu bytes.

  ER_BINLOG_CRASHED_BINLOG_TRIMMED was added in 8.0.2.

- Error: 4482 SQLSTATE: HY000 (ER_BINLOG_CANT_CLEAR_IN_USE_FLAG_FOR_CRASHED_BINLOG)

  Message: Failed to clear LOG_EVENT_BINLOG_IN_USE_F for the crashed binlog file when master server is recovering it.

  ER_BINLOG_CANT_CLEAR_IN_USE_FLAG_FOR_CRASHED_BINLOG was added in 8.0.2.

- Error: 4483 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_RUN_AFTER_SYNC_HOOK)

  Message: Failed to run 'after_sync' hooks

  ER_BINLOG_FAILED_TO_RUN_AFTER_SYNC_HOOK was added in 8.0.2.

- Error: 4484 SQLSTATE: HY000 (ER_TURNING_LOGGING_OFF_FOR_THE_DURATION)

  Message: %s Hence turning logging off for the whole duration of the MySQL server process. To turn it on again: fix the cause, shutdown the MySQL server and restart it.

  ER_TURNING_LOGGING_OFF_FOR_THE_DURATION was added in 8.0.2.

- Error: 4485 SQLSTATE: HY000 (ER_BINLOG_FAILED_TO_RUN_AFTER_FLUSH_HOOK)

  Message: Failed to run 'after_flush' hooks

  ER_BINLOG_FAILED_TO_RUN_AFTER_FLUSH_HOOK was added in 8.0.2.

- Error: 4486 SQLSTATE: HY000 (ER_BINLOG_CRASH_RECOVERY_FAILED)

  Message: Crash recovery failed. Either correct the problem (if it's, for example, out of memory error) and restart, or delete (or rename) binary log and start mysqld with --tc-heuristic-recover={commit|rollback}

  ER_BINLOG_CRASH_RECOVERY_FAILED was added in 8.0.2.

- Error: 4487 SQLSTATE: HY000 (ER_BINLOG_WARNING_SUPPRESSED)

Message: The following warning was suppressed %d times during the last %d seconds in the error log

`ER_BINLOG_WARNING_SUPPRESSED` was added in 8.0.2.

- Error: `4488` SQLSTATE: `HY000` (`ER_NDB_LOG_ENTRY`)

  Message: NDB: %s

  `ER_NDB_LOG_ENTRY` was added in 8.0.2.

- Error: `4489` SQLSTATE: `HY000` (`ER_NDB_LOG_ENTRY_WITH_PREFIX`)

  Message: NDB %s: %s

  `ER_NDB_LOG_ENTRY_WITH_PREFIX` was added in 8.0.2.

- Error: `4490` SQLSTATE: `HY000` (`ER_NDB_BINLOG_CANT_CREATE_PURGE_THD`)

  Message: NDB: Unable to purge %s.%s File=%s (failed to setup thd)

  `ER_NDB_BINLOG_CANT_CREATE_PURGE_THD` was added in 8.0.2.

- Error: `4491` SQLSTATE: `HY000` (`ER_INNODB_UNKNOWN_COLLATION`)

  Message: Unknown collation #%lu.

  `ER_INNODB_UNKNOWN_COLLATION` was added in 8.0.2.

- Error: `4492` SQLSTATE: `HY000` (`ER_INNODB_INVALID_LOG_GROUP_HOME_DIR`)

  Message: syntax error in innodb_log_group_home_dir

  `ER_INNODB_INVALID_LOG_GROUP_HOME_DIR` was added in 8.0.2.

- Error: `4493` SQLSTATE: `HY000` (`ER_INNODB_INVALID_INNODB_UNDO_DIRECTORY`)

  Message: syntax error in innodb_undo_directory

  `ER_INNODB_INVALID_INNODB_UNDO_DIRECTORY` was added in 8.0.2.

- Error: `4494` SQLSTATE: `HY000` (`ER_INNODB_ILLEGAL_COLON_IN_POOL`)

  Message: InnoDB: innodb_buffer_pool_filename cannot have colon (:) in the file name.

  `ER_INNODB_ILLEGAL_COLON_IN_POOL` was added in 8.0.2.

- Error: `4495` SQLSTATE: `HY000` (`ER_INNODB_INVALID_PAGE_SIZE`)

  Message: InnoDB: Invalid page size=%lu.

  `ER_INNODB_INVALID_PAGE_SIZE` was added in 8.0.2.

- Error: `4496` SQLSTATE: `HY000` (`ER_INNODB_DIRTY_WATER_MARK_NOT_LOW`)

  Message: InnoDB: innodb_max_dirty_pages_pct_lwm cannot be set higher than innodb_max_dirty_pages_pct. Setting innodb_max_dirty_pages_pct_lwm to %lf

  `ER_INNODB_DIRTY_WATER_MARK_NOT_LOW` was added in 8.0.2.

- Error: `4497` SQLSTATE: `HY000` (`ER_INNODB_IO_CAPACITY_EXCEEDS_MAX`)

Message: InnoDB: innodb_io_capacity cannot be set higher than innodb_io_capacity_max. Setting innodb_io_capacity to %lu

`ER_INNODB_IO_CAPACITY_EXCEEDS_MAX` was added in 8.0.2.

- Error: `4498` SQLSTATE: `HY000` (`ER_INNODB_FILES_SAME`)

  Message: %s and %s file names seem to be the same.

  `ER_INNODB_FILES_SAME` was added in 8.0.2.

- Error: `4499` SQLSTATE: `HY000` (`ER_INNODB_UNREGISTERED_TRX_ACTIVE`)

  Message: Transaction not registered for MySQL 2PC, but transaction is active

  `ER_INNODB_UNREGISTERED_TRX_ACTIVE` was added in 8.0.2.

- Error: `4500` SQLSTATE: `HY000` (`ER_INNODB_CLOSING_CONNECTION_ROLLS_BACK`)

  Message: MySQL is closing a connection that has an active InnoDB transaction. %s row modifications will roll back.

  `ER_INNODB_CLOSING_CONNECTION_ROLLS_BACK` was added in 8.0.2.

- Error: `4501` SQLSTATE: `HY000` (`ER_INNODB_TRX_XLATION_TABLE_OOM`)

  Message: InnoDB: fail to allocate memory for index translation table. Number of Index:%lu, array size:%lu

  `ER_INNODB_TRX_XLATION_TABLE_OOM` was added in 8.0.2.

- Error: `4502` SQLSTATE: `HY000` (`ER_INNODB_CANT_FIND_INDEX_IN_INNODB_DD`)

  Message: Cannot find index %s in InnoDB index dictionary.

  `ER_INNODB_CANT_FIND_INDEX_IN_INNODB_DD` was added in 8.0.2.

- Error: `4503` SQLSTATE: `HY000` (`ER_INNODB_INDEX_COLUMN_INFO_UNLIKE_MYSQLS`)

  Message: Found index %s whose column info does not match that of MySQL.

  `ER_INNODB_INDEX_COLUMN_INFO_UNLIKE_MYSQLS` was added in 8.0.2.

- Error: `4504` SQLSTATE: `HY000` (`ER_INNODB_CANT_OPEN_TABLE`)

  Message: Failed to open table %s.

  `ER_INNODB_CANT_OPEN_TABLE` was added in 8.0.2.

- Error: `4505` SQLSTATE: `HY000` (`ER_INNODB_CANT_BUILD_INDEX_XLATION_TABLE_FOR`)

  Message: Build InnoDB index translation table for Table %s failed

  `ER_INNODB_CANT_BUILD_INDEX_XLATION_TABLE_FOR` was added in 8.0.2.

- Error: `4506` SQLSTATE: `HY000` (`ER_INNODB_PK_NOT_IN_MYSQL`)

  Message: Table %s has a primary key in InnoDB data dictionary, but not in MySQL!

  `ER_INNODB_PK_NOT_IN_MYSQL` was added in 8.0.2.

- Error: `4507` SQLSTATE: `HY000` (`ER_INNODB_PK_ONLY_IN_MYSQL`)

Message: Table %s has no primary key in InnoDB data dictionary, but has one in MySQL! If you created the table with a MySQL version < 3.23.54 and did not define a primary key, but defined a unique key with all non-NULL columns, then MySQL internally treats that key as the primary key. You can fix this error by dump + DROP + CREATE + reimport of the table.

`ER_INNODB_PK_ONLY_IN_MYSQL` was added in 8.0.2.

- Error: `4508` SQLSTATE: `HY000` (`ER_INNODB_CLUSTERED_INDEX_PRIVATE`)

  Message: Table %s key_used_on_scan is %lu even though there is no primary key inside InnoDB.

  `ER_INNODB_CLUSTERED_INDEX_PRIVATE` was added in 8.0.2.

- Error: `4509` SQLSTATE: `HY000` (`ER_INNODB_PARTITION_TABLE_LOWERCASED`)

  Message: Partition table %s opened after converting to lower case. The table may have been moved from a case in-sensitive file system. Please recreate table in the current file system

  `ER_INNODB_PARTITION_TABLE_LOWERCASED` was added in 8.0.2.

- Error: `4510` SQLSTATE: `HY000` (`ER_ERRMSG_REPLACEMENT_DODGY`)

  Message: Cannot replace error message (%s,%s,%s) "%s" with "%s"; wrong number or type of %% subsitutions.

  `ER_ERRMSG_REPLACEMENT_DODGY` was added in 8.0.2.

- Error: `4511` SQLSTATE: `HY000` (`ER_ERRMSG_REPLACEMENTS_FAILED`)

  Message: Table for error message replacements could not be found or read, or one or more replacements could not be applied.

  `ER_ERRMSG_REPLACEMENTS_FAILED` was added in 8.0.2.

- Error: `4512` SQLSTATE: `HY000` (`ER_NPIPE_CANT_CREATE`)

  Message: %s: %s

  `ER_NPIPE_CANT_CREATE` was added in 8.0.2.

- Error: `4513` SQLSTATE: `HY000` (`ER_PARTITION_MOVE_CREATED_DUPLICATE_ROW_PLEASE_FIX`)

  Message: Table '%s': Delete from part %d failed with error %d. But it was already inserted into part %d, when moving the misplaced row! Please manually fix the duplicate row: %s

  `ER_PARTITION_MOVE_CREATED_DUPLICATE_ROW_PLEASE_FIX` was added in 8.0.2.

- Error: `4514` SQLSTATE: `HY000` (`ER_AUDIT_CANT_ABORT_COMMAND`)

  Message: Command '%s' cannot be aborted. The trigger error was (%d) [%s]: %s

  `ER_AUDIT_CANT_ABORT_COMMAND` was added in 8.0.2.

- Error: `4515` SQLSTATE: `HY000` (`ER_AUDIT_CANT_ABORT_EVENT`)

  Message: Event '%s' cannot be aborted. The trigger error was (%d) [%s]: %s

  `ER_AUDIT_CANT_ABORT_EVENT` was added in 8.0.2.

- Error: `4516` SQLSTATE: `HY000` (`ER_AUDIT_WARNING`)

  Message: %s. The trigger error was (%d) [%s]: %s

ER_AUDIT_WARNING was added in 8.0.2.

- Error: 4517 SQLSTATE: HY000 (ER_NDB_NUMBER_OF_CHANNELS)

  Message: Slave SQL: Configuration with number of replication masters = %u' is not supported when applying to Ndb

  ER_NDB_NUMBER_OF_CHANNELS was added in 8.0.2.

- Error: 4518 SQLSTATE: HY000 (ER_NDB_SLAVE_PARALLEL_WORKERS)

  Message: Slave SQL: Configuration 'slave_parallel_workers = %lu' is not supported when applying to Ndb

  ER_NDB_SLAVE_PARALLEL_WORKERS was added in 8.0.2.

- Error: 4519 SQLSTATE: HY000 (ER_NDB_DISTRIBUTING_ERR)

  Message: NDB %s: distributing %s err: %u

  ER_NDB_DISTRIBUTING_ERR was added in 8.0.2.

- Error: 4520 SQLSTATE: HY000 (ER_PART_EXPR_TOO_LONG)

  Message: Partitioning expression is too long.

  ER_PART_EXPR_TOO_LONG was added in 8.0.2.

- Error: 4521 SQLSTATE: HY000 (ER_UDF_DROP_DYNAMICALLY_REGISTERED)

  Message: DROP FUNCTION can't drop a dynamically registered user defined function

  ER_UDF_DROP_DYNAMICALLY_REGISTERED was added in 8.0.2.

- Error: 4522 SQLSTATE: HY000 (ER_UNABLE_TO_STORE_COLUMN_STATISTICS)

  Message: Unable to store column statistics for column '%s' in table '%s'.'%s'

  ER_UNABLE_TO_STORE_COLUMN_STATISTICS was added in 8.0.2.

- Error: 4523 SQLSTATE: HY000 (ER_UNABLE_TO_UPDATE_COLUMN_STATISTICS)

  Message: Unable to update column statistics for column '%s' in table '%s'.'%s'

  ER_UNABLE_TO_UPDATE_COLUMN_STATISTICS was added in 8.0.2.

- Error: 4524 SQLSTATE: HY000 (ER_UNABLE_TO_DROP_COLUMN_STATISTICS)

  Message: Unable to remove column statistics for column '%s' in table '%s'.'%s'

  ER_UNABLE_TO_DROP_COLUMN_STATISTICS was added in 8.0.2.

- Error: 4525 SQLSTATE: HY000 (ER_UNABLE_TO_BUILD_HISTOGRAM)

  Message: Unable to build histogram statistics for column '%s' in table '%s'.'%s'

  ER_UNABLE_TO_BUILD_HISTOGRAM was added in 8.0.2.

- Error: 4526 SQLSTATE: HY000 (ER_STOP_SLAVE_IO_THREAD_DISK_SPACE)

  Message: Waiting until I/O thread for channel '%s' finish writing to disk before stopping. Free some disk space or use 'KILL' to abort I/O thread operation. Notice that aborting the I/O thread while rotating the relay log might corrupt the relay logs, requiring a server restart to fix it.

ER_STOP_SLAVE_IO_THREAD_DISK_SPACE was added in 8.0.2.

- Error: 4527 SQLSTATE: HY000 (ER_MANDATORY_ROLE)

  Message: The role %s is a mandatory role and can't be revoked or dropped. The restriction can be lifted by excluding the role identifier from the global variable mandatory_roles.

  ER_MANDATORY_ROLE was added in 8.0.2.

- Error: 4528 SQLSTATE: HY000 (ER_MISSING_TABLESPACE_FILE)

  Message: Tablespace '%s' does not have a file named '%s'

  ER_MISSING_TABLESPACE_FILE was added in 8.0.3.

- Error: 4529 SQLSTATE: 42000 (ER_PERSIST_ONLY_ACCESS_DENIED_ERROR)

  Message: Access denied; you need %s privileges for this operation

  ER_PERSIST_ONLY_ACCESS_DENIED_ERROR was added in 8.0.3.

- Error: 4530 SQLSTATE: HY000 (ER_CMD_NEED_SUPER)

  Message: You need the SUPER privilege for command '%s'

  ER_CMD_NEED_SUPER was added in 8.0.2.

- Error: 4531 SQLSTATE: HY000 (ER_PATH_IN_DATADIR)

  Message: Path is within the current data directory '%s'

  ER_PATH_IN_DATADIR was added in 8.0.2.

- Error: 4532 SQLSTATE: HY000 (ER_DDL_IN_PROGRESS)

  Message: Concurrent DDL is performed during the operation. Please try again.

  ER_DDL_IN_PROGRESS was added in 8.0.2.

- Error: 4533 SQLSTATE: HY000 (ER_TOO_MANY_CONCURRENT_CLONES)

  Message: Too many concurrent clone operations. Maximum allowed - %d.

  ER_TOO_MANY_CONCURRENT_CLONES was added in 8.0.2.

B.4 Client Error Codes and Messages

Client error information comes from the following source files:

- The Error values and the symbols in parentheses correspond to definitions in the include/errmsg.h MySQL source file.
- The Message values correspond to the error messages that are listed in the libmysql/errmsg.c file. %d and %s represent numbers and strings, respectively, that are substituted into the messages when they are displayed.

Because updates are frequent, it is possible that those files will contain additional error information not listed here.

- Error: 2000 (CR_UNKNOWN_ERROR)

  Message: Unknown MySQL error

- Error: 2001 (CR_SOCKET_CREATE_ERROR)

  Message: Can't create UNIX socket (%d)

- Error: 2002 (CR_CONNECTION_ERROR)

  Message: Can't connect to local MySQL server through socket '%s' (%d)

- Error: 2003 (CR_CONN_HOST_ERROR)

  Message: Can't connect to MySQL server on '%s' (%d)

- Error: 2004 (CR_IPSOCK_ERROR)

  Message: Can't create TCP/IP socket (%d)

- Error: 2005 (CR_UNKNOWN_HOST)

  Message: Unknown MySQL server host '%s' (%d)

- Error: 2006 (CR_SERVER_GONE_ERROR)

  Message: MySQL server has gone away

- Error: 2007 (CR_VERSION_ERROR)

  Message: Protocol mismatch; server version = %d, client version = %d

- Error: 2008 (CR_OUT_OF_MEMORY)

  Message: MySQL client ran out of memory

- Error: 2009 (CR_WRONG_HOST_INFO)

  Message: Wrong host info

- Error: 2010 (CR_LOCALHOST_CONNECTION)

  Message: Localhost via UNIX socket

- Error: 2011 (CR_TCP_CONNECTION)

  Message: %s via TCP/IP

- Error: 2012 (CR_SERVER_HANDSHAKE_ERR)

  Message: Error in server handshake

- Error: 2013 (CR_SERVER_LOST)

  Message: Lost connection to MySQL server during query

- Error: 2014 (CR_COMMANDS_OUT_OF_SYNC)

  Message: Commands out of sync; you can't run this command now

- Error: 2015 (CR_NAMEDPIPE_CONNECTION)

  Message: Named pipe: %s

- Error: 2016 (CR_NAMEDPIPEWAIT_ERROR)

  Message: Can't wait for named pipe to host: %s pipe: %s (%lu)

- Error: 2017 (CR_NAMEDPIPEOPEN_ERROR)

Message: Can't open named pipe to host: %s pipe: %s (%lu)

- Error: 2018 (CR_NAMEDPIPESETSTATE_ERROR)

  Message: Can't set state of named pipe to host: %s pipe: %s (%lu)

- Error: 2019 (CR_CANT_READ_CHARSET)

  Message: Can't initialize character set %s (path: %s)

- Error: 2020 (CR_NET_PACKET_TOO_LARGE)

  Message: Got packet bigger than 'max_allowed_packet' bytes

- Error: 2021 (CR_EMBEDDED_CONNECTION)

  Message: Embedded server

- Error: 2022 (CR_PROBE_SLAVE_STATUS)

  Message: Error on SHOW SLAVE STATUS:

- Error: 2023 (CR_PROBE_SLAVE_HOSTS)

  Message: Error on SHOW SLAVE HOSTS:

- Error: 2024 (CR_PROBE_SLAVE_CONNECT)

  Message: Error connecting to slave:

- Error: 2025 (CR_PROBE_MASTER_CONNECT)

  Message: Error connecting to master:

- Error: 2026 (CR_SSL_CONNECTION_ERROR)

  Message: SSL connection error: %s

- Error: 2027 (CR_MALFORMED_PACKET)

  Message: Malformed packet

- Error: 2028 (CR_WRONG_LICENSE)

  Message: This client library is licensed only for use with MySQL servers having '%s' license

- Error: 2029 (CR_NULL_POINTER)

  Message: Invalid use of null pointer

- Error: 2030 (CR_NO_PREPARE_STMT)

  Message: Statement not prepared

- Error: 2031 (CR_PARAMS_NOT_BOUND)

  Message: No data supplied for parameters in prepared statement

- Error: 2032 (CR_DATA_TRUNCATED)

  Message: Data truncated

- Error: 2033 (CR_NO_PARAMETERS_EXISTS)

  Message: No parameters exist in the statement

- Error: 2034 (CR_INVALID_PARAMETER_NO)

  Message: Invalid parameter number

  The column number for `mysql_stmt_fetch_column()` was invalid.

  The parameter number for `mysql_stmt_send_long_data()` was invalid.

  A key name was empty or the amount of connection attribute data for `mysql_options4()` exceeds the 64KB limit.

- Error: 2035 (CR_INVALID_BUFFER_USE)

  Message: Can't send long data for non-string/non-binary data types (parameter: %d)

- Error: 2036 (CR_UNSUPPORTED_PARAM_TYPE)

  Message: Using unsupported buffer type: %d (parameter: %d)

- Error: 2037 (CR_SHARED_MEMORY_CONNECTION)

  Message: Shared memory: %s

- Error: 2038 (CR_SHARED_MEMORY_CONNECT_REQUEST_ERROR)

  Message: Can't open shared memory; client could not create request event (%lu)

- Error: 2039 (CR_SHARED_MEMORY_CONNECT_ANSWER_ERROR)

  Message: Can't open shared memory; no answer event received from server (%lu)

- Error: 2040 (CR_SHARED_MEMORY_CONNECT_FILE_MAP_ERROR)

  Message: Can't open shared memory; server could not allocate file mapping (%lu)

- Error: 2041 (CR_SHARED_MEMORY_CONNECT_MAP_ERROR)

  Message: Can't open shared memory; server could not get pointer to file mapping (%lu)

- Error: 2042 (CR_SHARED_MEMORY_FILE_MAP_ERROR)

  Message: Can't open shared memory; client could not allocate file mapping (%lu)

- Error: 2043 (CR_SHARED_MEMORY_MAP_ERROR)

  Message: Can't open shared memory; client could not get pointer to file mapping (%lu)

- Error: 2044 (CR_SHARED_MEMORY_EVENT_ERROR)

  Message: Can't open shared memory; client could not create %s event (%lu)

- Error: 2045 (CR_SHARED_MEMORY_CONNECT_ABANDONED_ERROR)

  Message: Can't open shared memory; no answer from server (%lu)

- Error: 2046 (CR_SHARED_MEMORY_CONNECT_SET_ERROR)

  Message: Can't open shared memory; cannot send request event to server (%lu)

- Error: 2047 (CR_CONN_UNKNOW_PROTOCOL)

  Message: Wrong or unknown protocol

- Error: 2048 (CR_INVALID_CONN_HANDLE)

Message: Invalid connection handle

- Error: `2049` (`CR_UNUSED_1`)

  Message: Connection using old (pre-4.1.1) authentication protocol refused (client option 'secure_auth' enabled)

- Error: `2050` (`CR_FETCH_CANCELED`)

  Message: Row retrieval was canceled by mysql_stmt_close() call

- Error: `2051` (`CR_NO_DATA`)

  Message: Attempt to read column without prior row fetch

- Error: `2052` (`CR_NO_STMT_METADATA`)

  Message: Prepared statement contains no metadata

- Error: `2053` (`CR_NO_RESULT_SET`)

  Message: Attempt to read a row while there is no result set associated with the statement

- Error: `2054` (`CR_NOT_IMPLEMENTED`)

  Message: This feature is not implemented yet

- Error: `2055` (`CR_SERVER_LOST_EXTENDED`)

  Message: Lost connection to MySQL server at '%s', system error: %d

- Error: `2056` (`CR_STMT_CLOSED`)

  Message: Statement closed indirectly because of a preceding %s() call

- Error: `2057` (`CR_NEW_STMT_METADATA`)

  Message: The number of columns in the result set differs from the number of bound buffers. You must reset the statement, rebind the result set columns, and execute the statement again

- Error: `2058` (`CR_ALREADY_CONNECTED`)

  Message: This handle is already connected. Use a separate handle for each connection.

- Error: `2059` (`CR_AUTH_PLUGIN_CANNOT_LOAD`)

  Message: Authentication plugin '%s' cannot be loaded: %s

- Error: `2060` (`CR_DUPLICATE_CONNECTION_ATTR`)

  Message: There is an attribute with the same name already

  A duplicate connection attribute name was specified for `mysql_options4()`.

- Error: `2061` (`CR_AUTH_PLUGIN_ERR`)

  Message: Authentication plugin '%s' reported error: %s

- Error: `2062` (`CR_INSECURE_API_ERR`)

  Message: Insecure API function call: '%s' Use instead: '%s'

  An insecure function call was detected. Modify the application to use the suggested alternative function instead.

- Error: 2063 (CR_FILE_NAME_TOO_LONG)

  Message: File name is too long

  CR_FILE_NAME_TOO_LONG was added in 8.0.1.

B.5 Problems and Common Errors

This section lists some common problems and error messages that you may encounter. It describes how to determine the causes of the problems and what to do to solve them.

B.5.1 How to Determine What Is Causing a Problem

When you run into a problem, the first thing you should do is to find out which program or piece of equipment is causing it:

- If you have one of the following symptoms, then it is probably a hardware problems (such as memory, motherboard, CPU, or hard disk) or kernel problem:
  - The keyboard does not work. This can normally be checked by pressing the Caps Lock key. If the Caps Lock light does not change, you have to replace your keyboard. (Before doing this, you should try to restart your computer and check all cables to the keyboard.)
  - The mouse pointer does not move.
  - The machine does not answer to a remote machine's pings.
  - Other programs that are not related to MySQL do not behave correctly.
  - Your system restarted unexpectedly. (A faulty user-level program should never be able to take down your system.)

  In this case, you should start by checking all your cables and run some diagnostic tool to check your hardware! You should also check whether there are any patches, updates, or service packs for your operating system that could likely solve your problem. Check also that all your libraries (such as glibc) are up to date.

  It is always good to use a machine with ECC memory to discover memory problems early.
- If your keyboard is locked up, you may be able to recover by logging in to your machine from another machine and executing kbd_mode -a.
- Please examine your system log file (/var/log/messages or similar) for reasons for your problem. If you think the problem is in MySQL, you should also examine MySQL's log files. See Section 5.4, “MySQL Server Logs”.
- If you do not think you have hardware problems, you should try to find out which program is causing problems. Try using top, ps, Task Manager, or some similar program, to check which program is taking all CPU or is locking the machine.
- Use top, df, or a similar program to check whether you are out of memory, disk space, file descriptors, or some other critical resource.
- If the problem is some runaway process, you can always try to kill it. If it does not want to die, there is probably a bug in the operating system.

If after you have examined all other possibilities and you have concluded that the MySQL server or a MySQL client is causing the problem, it is time to create a bug report for our mailing list or our support team. In the bug report, try to give a very detailed description of how the system is behaving and what you think is happening. You should also state why you think that MySQL is causing the problem. Take

into consideration all the situations in this chapter. State any problems exactly how they appear when you examine your system. Use the “copy and paste” method for any output and error messages from programs and log files.

Try to describe in detail which program is not working and all symptoms you see. We have in the past received many bug reports that state only “the system does not work.” This provides us with no information about what could be the problem.

If a program fails, it is always useful to know the following information:

- Has the program in question made a segmentation fault (did it dump core)?
- Is the program taking up all available CPU time? Check with `top`. Let the program run for a while, it may simply be evaluating something computationally intensive.
- If the `mysqld` server is causing problems, can you get any response from it with `mysqladmin -u root ping` or `mysqladmin -u root processlist`?
- What does a client program say when you try to connect to the MySQL server? (Try with `mysql`, for example.) Does the client jam? Do you get any output from the program?

When sending a bug report, you should follow the outline described in Section 1.7, “How to Report Bugs or Problems”.

B.5.2 Common Errors When Using MySQL Programs

This section lists some errors that users frequently encounter when running MySQL programs. Although the problems show up when you try to run client programs, the solutions to many of the problems involves changing the configuration of the MySQL server.

B.5.2.1 Access denied

An `Access denied` error can have many causes. Often the problem is related to the MySQL accounts that the server permits client programs to use when connecting. See Section 6.2, “The MySQL Access Privilege System”, and Section 6.2.9, “Troubleshooting Problems Connecting to MySQL”.

B.5.2.2 Can't connect to [local] MySQL server

A MySQL client on Unix can connect to the `mysqld` server in two different ways: By using a Unix socket file to connect through a file in the file system (default `/tmp/mysql.sock`), or by using TCP/IP, which connects through a port number. A Unix socket file connection is faster than TCP/IP, but can be used only when connecting to a server on the same computer. A Unix socket file is used if you do not specify a host name or if you specify the special host name `localhost`.

If the MySQL server is running on Windows, you can connect using TCP/IP. If the server is started with the `--enable-named-pipe` option, you can also connect with named pipes if you run the client on the host where the server is running. The name of the named pipe is `MySQL` by default. If you do not give a host name when connecting to `mysqld`, a MySQL client first tries to connect to the named pipe. If that does not work, it connects to the TCP/IP port. You can force the use of named pipes on Windows by using `.` as the host name.

The error (2002) `Can't connect to ...` normally means that there is no MySQL server running on the system or that you are using an incorrect Unix socket file name or TCP/IP port number when trying to connect to the server. You should also check that the TCP/IP port you are using has not been blocked by a firewall or port blocking service.

The error (2003) `Can't connect to MySQL server on 'server' (10061)` indicates that the network connection has been refused. You should check that there is a MySQL server running, that it

has network connections enabled, and that the network port you specified is the one configured on the server.

Start by checking whether there is a process named `mysqld` running on your server host. (Use `ps xa | grep mysqld` on Unix or the Task Manager on Windows.) If there is no such process, you should start the server. See Section 2.9.2, “Starting the Server”.

If a `mysqld` process is running, you can check it by trying the following commands. The port number or Unix socket file name might be different in your setup. `host_ip` represents the IP address of the machine where the server is running.

```
shell> mysqladmin version
shell> mysqladmin variables
shell> mysqladmin -h `hostname` version variables
shell> mysqladmin -h `hostname` --port=3306 version
shell> mysqladmin -h host_ip version
shell> mysqladmin --protocol=SOCKET --socket=/tmp/mysql.sock version
```

Note the use of backticks rather than forward quotation marks with the `hostname` command; these cause the output of `hostname` (that is, the current host name) to be substituted into the `mysqladmin` command. If you have no `hostname` command or are running on Windows, you can manually type the host name of your machine (without backticks) following the `-h` option. You can also try `-h 127.0.0.1` to connect with TCP/IP to the local host.

Make sure that the server has not been configured to ignore network connections or (if you are attempting to connect remotely) that it has not been configured to listen only locally on its network interfaces. If the server was started with `--skip-networking`, it will not accept TCP/IP connections at all. If the server was started with `--bind-address=127.0.0.1`, it will listen for TCP/IP connections only locally on the loopback interface and will not accept remote connections.

Check to make sure that there is no firewall blocking access to MySQL. Your firewall may be configured on the basis of the application being executed, or the port number used by MySQL for communication (3306 by default). Under Linux or Unix, check your IP tables (or similar) configuration to ensure that the port has not been blocked. Under Windows, applications such as ZoneAlarm or Windows Firewall may need to be configured not to block the MySQL port.

Here are some reasons the `Can't connect to local MySQL server` error might occur:

- `mysqld` is not running on the local host. Check your operating system's process list to ensure the `mysqld` process is present.
- You're running a MySQL server on Windows with many TCP/IP connections to it. If you're experiencing that quite often your clients get that error, you can find a workaround here: Connection to MySQL Server Failing on Windows.
- Someone has removed the Unix socket file that `mysqld` uses (`/tmp/mysql.sock` by default). For example, you might have a `cron` job that removes old files from the `/tmp` directory. You can always run `mysqladmin version` to check whether the Unix socket file that `mysqladmin` is trying to use really exists. The fix in this case is to change the `cron` job to not remove `mysql.sock` or to place the socket file somewhere else. See Section B.5.3.6, “How to Protect or Change the MySQL Unix Socket File”.
- You have started the `mysqld` server with the `--socket=/path/to/socket` option, but forgotten to tell client programs the new name of the socket file. If you change the socket path name for the server, you must also notify the MySQL clients. You can do this by providing the same `--socket` option when you run client programs. You also need to ensure that clients have permission to access the `mysql.sock` file. To find out where the socket file is, you can do:

  ```
  shell> netstat -ln | grep mysql
  ```

  See Section B.5.3.6, “How to Protect or Change the MySQL Unix Socket File”.

- You are using Linux and one server thread has died (dumped core). In this case, you must kill the other `mysqld` threads (for example, with `kill`) before you can restart the MySQL server. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.
- The server or client program might not have the proper access privileges for the directory that holds the Unix socket file or the socket file itself. In this case, you must either change the access privileges for the directory or socket file so that the server and clients can access them, or restart `mysqld` with a `--socket` option that specifies a socket file name in a directory where the server can create it and where client programs can access it.

If you get the error message `Can't connect to MySQL server on some_host`, you can try the following things to find out what the problem is:

- Check whether the server is running on that host by executing `telnet some_host 3306` and pressing the Enter key a couple of times. (3306 is the default MySQL port number. Change the value if your server is listening to a different port.) If there is a MySQL server running and listening to the port, you should get a response that includes the server's version number. If you get an error such as `telnet: Unable to connect to remote host: Connection refused`, then there is no server running on the given port.
- If the server is running on the local host, try using `mysqladmin -h localhost variables` to connect using the Unix socket file. Verify the TCP/IP port number that the server is configured to listen to (it is the value of the `port` variable.)
- If you are running under Linux and Security-Enhanced Linux (SELinux) is enabled, make sure you have disabled SELinux protection for the `mysqld` process.

Connection to MySQL Server Failing on Windows

When you're running a MySQL server on Windows with many TCP/IP connections to it, and you're experiencing that quite often your clients get a `Can't connect to MySQL server` error, the reason might be that Windows does not allow for enough ephemeral (short-lived) ports to serve those connections.

The purpose of `TIME_WAIT` is to keep a connection accepting packets even after the connection has been closed. This is because Internet routing can cause a packet to take a slow route to its destination and it may arrive after both sides have agreed to close. If the port is in use for a new connection, that packet from the old connection could break the protocol or compromise personal information from the original connection. The `TIME_WAIT` delay prevents this by ensuring that the port cannot be reused until after some time has been permitted for those delayed packets to arrive.

It is safe to reduce `TIME_WAIT` greatly on LAN connections because there is little chance of packets arriving at very long delays, as they could through the Internet with its comparatively large distances and latencies.

Windows permits ephemeral (short-lived) TCP ports to the user. After any port is closed it will remain in a `TIME_WAIT` status for 120 seconds. The port will not be available again until this time expires. The default range of port numbers depends on the version of Windows, with a more limited number of ports in older versions:

- Windows through Server 2003: Ports in range 1025–5000
- Windows Vista, Server 2008, and newer: Ports in range 49152–65535

With a small stack of available TCP ports (5000) and a high number of TCP ports being open and closed over a short period of time along with the `TIME_WAIT` status you have a good chance for running out of ports. There are two ways to address this problem:

- Reduce the number of TCP ports consumed quickly by investigating connection pooling or persistent connections where possible
- Tune some settings in the Windows registry (see below)

Important

The following procedure involves modifying the Windows registry. Before you modify the registry, make sure to back it up and make sure that you understand how to restore it if a problem occurs. For information about how to back up, restore, and edit the registry, view the following article in the Microsoft Knowledge Base: http://support.microsoft.com/kb/256986/EN-US/.

1. Start Registry Editor (`Regedt32.exe`).
2. Locate the following key in the registry:

   ```
   HKEY_LOCAL_MACHINE\SYSTEM\CurrentControlSet\Services\Tcpip\Parameters
   ```

3. On the `Edit` menu, click `Add Value`, and then add the following registry value:

   ```
   Value Name: MaxUserPort
   Data Type: REG_DWORD
   Value: 65534
   ```

   This sets the number of ephemeral ports available to any user. The valid range is between 5000 and 65534 (decimal). The default value is 0x1388 (5000 decimal).

4. On the `Edit` menu, click `Add Value`, and then add the following registry value:

   ```
   Value Name: TcpTimedWaitDelay
   Data Type: REG_DWORD
   Value: 30
   ```

   This sets the number of seconds to hold a TCP port connection in `TIME_WAIT` state before closing. The valid range is between 30 and 300 decimal, although you may wish to check with Microsoft for the latest permitted values. The default value is 0x78 (120 decimal).

5. Quit Registry Editor.
6. Reboot the machine.

Note: Undoing the above should be as simple as deleting the registry entries you've created.

B.5.2.3 Lost connection to MySQL server

There are three likely causes for this error message.

Usually it indicates network connectivity trouble and you should check the condition of your network if this error occurs frequently. If the error message includes “during query,” this is probably the case you are experiencing.

Sometimes the “during query” form happens when millions of rows are being sent as part of one or more queries. If you know that this is happening, you should try increasing `net_read_timeout` from its default of 30 seconds to 60 seconds or longer, sufficient for the data transfer to complete.

More rarely, it can happen when the client is attempting the initial connection to the server. In this case, if your `connect_timeout` value is set to only a few seconds, you may be able to resolve the problem by increasing it to ten seconds, perhaps more if you have a very long distance or slow connection. You can determine whether you are experiencing this more uncommon cause by using `SHOW GLOBAL STATUS LIKE 'Aborted_connects'`. It will increase by one for each initial connection attempt that the server aborts. You may see “reading authorization packet” as part of the error message; if so, that also suggests that this is the solution that you need.

If the cause is none of those just described, you may be experiencing a problem with `BLOB` values that are larger than `max_allowed_packet`, which can cause this error with some clients. Sometime

you may see an ER_NET_PACKET_TOO_LARGE error, and that confirms that you need to increase max_allowed_packet.

B.5.2.4 Password Fails When Entered Interactively

MySQL client programs prompt for a password when invoked with a --password or -p option that has no following password value:

```
shell> mysql -u user_name -p
Enter password:
```

On some systems, you may find that your password works when specified in an option file or on the command line, but not when you enter it interactively at the Enter password: prompt. This occurs when the library provided by the system to read passwords limits password values to a small number of characters (typically eight). That is a problem with the system library, not with MySQL. To work around it, change your MySQL password to a value that is eight or fewer characters long, or put your password in an option file.

B.5.2.5 Host 'host_name' is blocked

If the following error occurs, it means that mysqld has received many connection requests from the given host that were interrupted in the middle:

```
Host 'host_name' is blocked because of many connection errors.
Unblock with 'mysqladmin flush-hosts'
```

The value of the max_connect_errors system variable determines how many successive interrupted connection requests are permitted. (See Section 5.1.5, “Server System Variables”.) After max_connect_errors failed requests without a successful connection, mysqld assumes that something is wrong (for example, that someone is trying to break in), and blocks the host from further connections until you issue a FLUSH HOSTS statement or execute a mysqladmin flush-hosts command.

By default, mysqld blocks a host after 100 connection errors. You can adjust the value by setting max_connect_errors at server startup:

```
shell> mysqld_safe --max_connect_errors=10000 &
```

The value can also be set at runtime:

```
mysql> SET GLOBAL max_connect_errors=10000;
```

If you get the Host 'host_name' is blocked error message for a given host, you should first verify that there is nothing wrong with TCP/IP connections from that host. If you are having network problems, it does you no good to increase the value of the max_connect_errors variable.

B.5.2.6 Too many connections

If you get a Too many connections error when you try to connect to the mysqld server, this means that all available connections are in use by other clients.

The number of connections permitted is controlled by the max_connections system variable. The default value is 151 to improve performance when MySQL is used with the Apache Web server. (Previously, the default was 100.) If you need to support more connections, you should set a larger value for this variable.

mysqld actually permits max_connections+1 clients to connect. The extra connection is reserved for use by accounts that have the SUPER privilege. By granting the SUPER privilege to administrators and not to normal users (who should not need it), an administrator can connect to the server and use

`SHOW PROCESSLIST` to diagnose problems even if the maximum number of unprivileged clients are connected. See Section 13.7.5.29, “SHOW PROCESSLIST Syntax”.

The maximum number of connections MySQL supports depends on the quality of the thread library on a given platform, the amount of RAM available, how much RAM is used for each connection, the workload from each connection, and the desired response time. Linux or Solaris should be able to support at least 500 to 1000 simultaneous connections routinely and as many as 10,000 connections if you have many gigabytes of RAM available and the workload from each is low or the response time target undemanding. Windows is limited to (open tables × 2 + open connections) < 2048 due to the Posix compatibility layer used on that platform.

Increasing `open-files-limit` may be necessary. Also see Section 2.5, “Installing MySQL on Linux”, for how to raise the operating system limit on how many handles can be used by MySQL.

B.5.2.7 Out of memory

If you issue a query using the `mysql` client program and receive an error like the following one, it means that `mysql` does not have enough memory to store the entire query result:

```
mysql: Out of memory at line 42, 'malloc.c'
mysql: needed 8136 byte (8k), memory in use: 12481367 bytes (12189k)
ERROR 2008: MySQL client ran out of memory
```

To remedy the problem, first check whether your query is correct. Is it reasonable that it should return so many rows? If not, correct the query and try again. Otherwise, you can invoke `mysql` with the `--quick` option. This causes it to use the `mysql_use_result()` C API function to retrieve the result set, which places less of a load on the client (but more on the server).

B.5.2.8 MySQL server has gone away

This section also covers the related `Lost connection to server during query` error.

The most common reason for the `MySQL server has gone away` error is that the server timed out and closed the connection. In this case, you normally get one of the following error codes (which one you get is operating system-dependent).

Error Code	Description
`CR_SERVER_GONE_ERROR`	The client couldn't send a question to the server.
`CR_SERVER_LOST`	The client didn't get an error when writing to the server, but it didn't get a full answer (or any answer) to the question.

By default, the server closes the connection after eight hours if nothing has happened. You can change the time limit by setting the `wait_timeout` variable when you start `mysqld`. See Section 5.1.5, “Server System Variables”.

If you have a script, you just have to issue the query again for the client to do an automatic reconnection. This assumes that you have automatic reconnection in the client enabled (which is the default for the `mysql` command-line client).

Some other common reasons for the `MySQL server has gone away` error are:

- You (or the db administrator) has killed the running thread with a `KILL` statement or a `mysqladmin kill` command.
- You tried to run a query after closing the connection to the server. This indicates a logic error in the application that should be corrected.
- A client application running on a different host does not have the necessary privileges to connect to the MySQL server from that host.

- You got a timeout from the TCP/IP connection on the client side. This may happen if you have been using the commands: `mysql_options(..., MYSQL_OPT_READ_TIMEOUT,...)` or `mysql_options(..., MYSQL_OPT_WRITE_TIMEOUT,...)`. In this case increasing the timeout may help solve the problem.
- You have encountered a timeout on the server side and the automatic reconnection in the client is disabled (the `reconnect` flag in the `MYSQL` structure is equal to 0).
- You are using a Windows client and the server had dropped the connection (probably because `wait_timeout` expired) before the command was issued.

  The problem on Windows is that in some cases MySQL does not get an error from the OS when writing to the TCP/IP connection to the server, but instead gets the error when trying to read the answer from the connection.

  The solution to this is to either do a `mysql_ping()` on the connection if there has been a long time since the last query (this is what Connector/ODBC does) or set `wait_timeout` on the `mysqld` server so high that it in practice never times out.
- You can also get these errors if you send a query to the server that is incorrect or too large. If `mysqld` receives a packet that is too large or out of order, it assumes that something has gone wrong with the client and closes the connection. If you need big queries (for example, if you are working with big `BLOB` columns), you can increase the query limit by setting the server's `max_allowed_packet` variable, which has a default value of 4MB. You may also need to increase the maximum packet size on the client end. More information on setting the packet size is given in Section B.5.2.9, “Packet Too Large”.

  An `INSERT` or `REPLACE` statement that inserts a great many rows can also cause these sorts of errors. Either one of these statements sends a single request to the server irrespective of the number of rows to be inserted; thus, you can often avoid the error by reducing the number of rows sent per `INSERT` or `REPLACE`.
- You also get a lost connection if you are sending a packet 16MB or larger if your client is older than 4.0.8 and your server is 4.0.8 and above, or the other way around.
- It is also possible to see this error if host name lookups fail (for example, if the DNS server on which your server or network relies goes down). This is because MySQL is dependent on the host system for name resolution, but has no way of knowing whether it is working—from MySQL's point of view the problem is indistinguishable from any other network timeout.

  You may also see the `MySQL server has gone away` error if MySQL is started with the `--skip-networking` option.

  Another networking issue that can cause this error occurs if the MySQL port (default 3306) is blocked by your firewall, thus preventing any connections at all to the MySQL server.
- You can also encounter this error with applications that fork child processes, all of which try to use the same connection to the MySQL server. This can be avoided by using a separate connection for each child process.
- You have encountered a bug where the server died while executing the query.

You can check whether the MySQL server died and restarted by executing `mysqladmin version` and examining the server's uptime. If the client connection was broken because `mysqld` crashed and restarted, you should concentrate on finding the reason for the crash. Start by checking whether issuing the query again kills the server again. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.

You can get more information about the lost connections by starting `mysqld` with the `log_error_verbosity` system variable set to 3. This logs some of the disconnection messages in the `hostname.err` file. See Section 5.4.2, “The Error Log”.

If you want to create a bug report regarding this problem, be sure that you include the following information:

- Indicate whether the MySQL server died. You can find information about this in the server error log. See Section B.5.3.3, “What to Do If MySQL Keeps Crashing”.
- If a specific query kills `mysqld` and the tables involved were checked with `CHECK TABLE` before you ran the query, can you provide a reproducible test case? See Section 28.5, “Debugging and Porting MySQL”.
- What is the value of the `wait_timeout` system variable in the MySQL server? (`mysqladmin variables` gives you the value of this variable.)
- Have you tried to run `mysqld` with the general query log enabled to determine whether the problem query appears in the log? (See Section 5.4.3, “The General Query Log”.)

See also Section B.5.2.10, “Communication Errors and Aborted Connections”, and Section 1.7, “How to Report Bugs or Problems”.

B.5.2.9 Packet Too Large

A communication packet is a single SQL statement sent to the MySQL server, a single row that is sent to the client, or a binary log event sent from a master replication server to a slave.

The largest possible packet that can be transmitted to or from a MySQL 8.0 server or client is 1GB.

When a MySQL client or the `mysqld` server receives a packet bigger than `max_allowed_packet` bytes, it issues an `ER_NET_PACKET_TOO_LARGE` error and closes the connection. With some clients, you may also get a `Lost connection to MySQL server during query` error if the communication packet is too large.

Both the client and the server have their own `max_allowed_packet` variable, so if you want to handle big packets, you must increase this variable both in the client and in the server.

If you are using the `mysql` client program, its default `max_allowed_packet` variable is 16MB. To set a larger value, start `mysql` like this:

```
shell> mysql --max_allowed_packet=32M
```

That sets the packet size to 32MB.

The server's default `max_allowed_packet` value is 4MB. You can increase this if the server needs to handle big queries (for example, if you are working with big `BLOB` columns). For example, to set the variable to 16MB, start the server like this:

```
shell> mysqld --max_allowed_packet=16M
```

You can also use an option file to set `max_allowed_packet`. For example, to set the size for the server to 16MB, add the following lines in an option file:

```
[mysqld]
max_allowed_packet=16M
```

It is safe to increase the value of this variable because the extra memory is allocated only when needed. For example, `mysqld` allocates more memory only when you issue a long query or when `mysqld` must return a large result row. The small default value of the variable is a precaution to catch incorrect packets between the client and server and also to ensure that you do not run out of memory by using large packets accidentally.

You can also get strange problems with large packets if you are using large `BLOB` values but have not given `mysqld` access to enough memory to handle the query. If you suspect this is the case, try adding `ulimit -d 256000` to the beginning of the `mysqld_safe` script and restarting `mysqld`.

B.5.2.10 Communication Errors and Aborted Connections

If connection problems occur such as communication errors or aborted connections, use these sources of information to diagnose problems:

- The error log. See Section 5.4.2, “The Error Log”.
- The general query log. See Section 5.4.3, “The General Query Log”.
- The `Aborted_xxx` and `Connection_errors_xxx` status variables. See Section 5.1.7, “Server Status Variables”.
- The host cache, which is accessible using the `host_cache` Performance Schema table. See Section 8.12.4.2, “DNS Lookup Optimization and the Host Cache”, and Section 25.11.16.1, “The host_cache Table”.

If you start the server with the `log_error_verbosity` system variable set to 3, you might find messages like this in your error log:

```
[Note] Aborted connection 854 to db: 'employees' user: 'josh'
```

If a client is unable even to connect, the server increments the `Aborted_connects` status variable. Unsuccessful connection attempts can occur for the following reasons:

- A client attempts to access a database but has no privileges for it.
- A client uses an incorrect password.
- A connection packet does not contain the right information.
- It takes more than `connect_timeout` seconds to obtain a connect packet. See Section 5.1.5, “Server System Variables”.

If these kinds of things happen, it might indicate that someone is trying to break into your server! If the general query log is enabled, messages for these types of problems are logged to it.

If a client successfully connects but later disconnects improperly or is terminated, the server increments the `Aborted_clients` status variable, and logs an `Aborted connection` message to the error log. The cause can be any of the following:

- The client program did not call `mysql_close()` before exiting.
- The client had been sleeping more than `wait_timeout` or `interactive_timeout` seconds without issuing any requests to the server. See Section 5.1.5, “Server System Variables”.
- The client program ended abruptly in the middle of a data transfer.

Other reasons for problems with aborted connections or aborted clients:

- The `max_allowed_packet` variable value is too small or queries require more memory than you have allocated for `mysqld`. See Section B.5.2.9, “Packet Too Large”.
- Use of Ethernet protocol with Linux, both half and full duplex. Some Linux Ethernet drivers have this bug. You should test for this bug by transferring a huge file using FTP between the client and server machines. If a transfer goes in burst-pause-burst-pause mode, you are experiencing a Linux duplex syndrome. Switch the duplex mode for both your network card and hub/switch to either full duplex or to half duplex and test the results to determine the best setting.

- A problem with the thread library that causes interrupts on reads.
- Badly configured TCP/IP.
- Faulty Ethernets, hubs, switches, cables, and so forth. This can be diagnosed properly only by replacing hardware.

See also Section B.5.2.8, “MySQL server has gone away”.

B.5.2.11 The table is full

If a table-full error occurs, it may be that the disk is full or that the table has reached its maximum size. The effective maximum table size for MySQL databases is usually determined by operating system constraints on file sizes, not by MySQL internal limits. See Section C.10.3, “Limits on Table Size”.

B.5.2.12 Can't create/write to file

If you get an error of the following type for some queries, it means that MySQL cannot create a temporary file for the result set in the temporary directory:

```
Can't create/write to file '\\sqla3fe_0.ism'.
```

The preceding error is a typical message for Windows; the Unix message is similar.

One fix is to start `mysqld` with the `--tmpdir` option or to add the option to the `[mysqld]` section of your option file. For example, to specify a directory of `C:\temp`, use these lines:

```
[mysqld]
tmpdir=C:/temp
```

The `C:\temp` directory must exist and have sufficient space for the MySQL server to write to. See Section 4.2.6, “Using Option Files”.

Another cause of this error can be permissions issues. Make sure that the MySQL server can write to the `tmpdir` directory.

Check also the error code that you get with `perror`. One reason the server cannot write to a table is that the file system is full:

```
shell> perror 28
OS error code  28:  No space left on device
```

If you get an error of the following type during startup, it indicates that the file system or directory used for storing data files is write protected. Provided that the write error is to a test file, the error is not serious and can be safely ignored.

```
Can't create test file /usr/local/mysql/data/master.lower-test
```

B.5.2.13 Commands out of sync

If you get `Commands out of sync; you can't run this command now` in your client code, you are calling client functions in the wrong order.

This can happen, for example, if you are using `mysql_use_result()` and try to execute a new query before you have called `mysql_free_result()`. It can also happen if you try to execute two queries that return data without calling `mysql_use_result()` or `mysql_store_result()` in between.

B.5.2.14 Ignoring user

If you get the following error, it means that when `mysqld` was started or when it reloaded the grant tables, it found an account in the `user` table that had an invalid password.

```
Found wrong password for user 'some_user'@'some_host'; ignoring user
```

As a result, the account is simply ignored by the permission system. To fix this problem, assign a new, valid password to the account.

B.5.2.15 Table 'tbl_name' doesn't exist

If you get either of the following errors, it usually means that no table exists in the default database with the given name:

```
Table 'tbl_name' doesn't exist
Can't find file: 'tbl_name' (errno: 2)
```

In some cases, it may be that the table does exist but that you are referring to it incorrectly:

- Because MySQL uses directories and files to store databases and tables, database and table names are case sensitive if they are located on a file system that has case-sensitive file names.
- Even for file systems that are not case sensitive, such as on Windows, all references to a given table within a query must use the same lettercase.

You can check which tables are in the default database with `SHOW TABLES`. See Section 13.7.5, “SHOW Syntax”.

B.5.2.16 Can't initialize character set

You might see an error like this if you have character set problems:

```
MySQL Connection Failed: Can't initialize character set charset_name
```

This error can have any of the following causes:

- The character set is a multibyte character set and you have no support for the character set in the client. In this case, you need to recompile the client by running `CMake` with the `-DDEFAULT_CHARSET=charset_name` option. See Section 2.8.4, “MySQL Source-Configuration Options”.

  All standard MySQL binaries are compiled with support for all multibyte character sets.

- The character set is a simple character set that is not compiled into `mysqld`, and the character set definition files are not in the place where the client expects to find them.

  In this case, you need to use one of the following methods to solve the problem:

  - Recompile the client with support for the character set. See Section 2.8.4, “MySQL Source-Configuration Options”.
  - Specify to the client the directory where the character set definition files are located. For many clients, you can do this with the `--character-sets-dir` option.
  - Copy the character definition files to the path where the client expects them to be.

B.5.2.17 File Not Found and Similar Errors

If you get `ERROR 'file_name' not found (errno: 23)`, `Can't open file: file_name (errno: 24)`, or any other error with `errno 23` or `errno 24` from MySQL, it means that you have not allocated enough file descriptors for the MySQL server. You can use the `perror` utility to get a description of what the error number means:

```
shell> perror 23
OS error code  23:  File table overflow
shell> perror 24
OS error code  24:  Too many open files
shell> perror 11
OS error code  11:  Resource temporarily unavailable
```

The problem here is that `mysqld` is trying to keep open too many files simultaneously. You can either tell `mysqld` not to open so many files at once or increase the number of file descriptors available to `mysqld`.

To tell `mysqld` to keep open fewer files at a time, you can make the table cache smaller by reducing the value of the `table_open_cache` system variable (the default value is 64). This may not entirely prevent running out of file descriptors because in some circumstances the server may attempt to extend the cache size temporarily, as described in Section 8.4.3.1, “How MySQL Opens and Closes Tables”. Reducing the value of `max_connections` also reduces the number of open files (the default value is 100).

To change the number of file descriptors available to `mysqld`, you can use the `--open-files-limit` option to `mysqld_safe` or set the `open_files_limit` system variable. See Section 5.1.5, “Server System Variables”. The easiest way to set these values is to add an option to your option file. See Section 4.2.6, “Using Option Files”. If you have an old version of `mysqld` that does not support setting the open files limit, you can edit the `mysqld_safe` script. There is a commented-out line `ulimit -n 256` in the script. You can remove the `#` character to uncomment this line, and change the number `256` to set the number of file descriptors to be made available to `mysqld`.

`--open-files-limit` and `ulimit` can increase the number of file descriptors, but only up to the limit imposed by the operating system. There is also a “hard” limit that can be overridden only if you start `mysqld_safe` or `mysqld` as `root` (just remember that you also need to start the server with the `--user` option in this case so that it does not continue to run as `root` after it starts up). If you need to increase the operating system limit on the number of file descriptors available to each process, consult the documentation for your system.

Note

If you run the `tcsh` shell, `ulimit` does not work! `tcsh` also reports incorrect values when you ask for the current limits. In this case, you should start `mysqld_safe` using `sh`.

B.5.2.18 Table-Corruption Issues

If you have started `mysqld` with `--myisam-recover-options`, MySQL automatically checks and tries to repair `MyISAM` tables if they are marked as 'not closed properly' or 'crashed'. If this happens, MySQL writes an entry in the `hostname.err` file `'Warning: Checking table ...'` which is followed by `Warning: Repairing table` if the table needs to be repaired. If you get a lot of these errors, without `mysqld` having died unexpectedly just before, then something is wrong and needs to be investigated further.

When the server detects `MyISAM` table corruption, it writes additional information to the error log, such as the name and line number of the source file, and the list of threads accessing the table. Example: `Got an error from thread_id=1, mi_dynrec.c:368`. This is useful information to include in bug reports.

See also Section 5.1.4, “Server Command Options”, and Section 28.5.1.7, “Making a Test Case If You Experience Table Corruption”.

B.5.3 Administration-Related Issues

B.5.3.1 Problems with File Permissions

If you have problems with file permissions, the `UMASK` or `UMASK_DIR` environment variable might be set incorrectly when `mysqld` starts. For example, MySQL might issue the following error message when you create a table:

```
ERROR: Can't find file: 'path/with/file_name' (Errcode: 13)
```

The default `UMASK` and `UMASK_DIR` values are `0640` and `0750`, respectively. MySQL assumes that the value for `UMASK` or `UMASK_DIR` is in octal if it starts with a zero. For example, setting `UMASK=0600` is equivalent to `UMASK=384` because 0600 octal is 384 decimal.

To change the default `UMASK` value, start `mysqld_safe` as follows:

```
shell> UMASK=384  # = 600 in octal
shell> export UMASK
shell> mysqld_safe &
```

By default, MySQL creates database directories with an access permission value of `0750`. To modify this behavior, set the `UMASK_DIR` variable. If you set its value, new directories are created with the combined `UMASK` and `UMASK_DIR` values. For example, to give group access to all new directories, start `mysqld_safe` as follows:

```
shell> UMASK_DIR=504  # = 770 in octal
shell> export UMASK_DIR
shell> mysqld_safe &
```

For additional details, see Section 4.9, “MySQL Program Environment Variables”.

B.5.3.2 How to Reset the Root Password

If you have never assigned a `root` password for MySQL, the server does not require a password at all for connecting as `root`. However, this is insecure. For instructions on assigning a password, see Section 2.9.4, “Securing the Initial MySQL Account”.

If you know the `root` password and want to change it, see Section 13.7.1.1, “ALTER USER Syntax”, and Section 13.7.1.10, “SET PASSWORD Syntax”.

If you assigned a `root` password previously but have forgotten it, you can assign a new password. The following sections provide instructions for Windows and Unix and Unix-like systems, as well as generic instructions that apply to any system.

Resetting the Root Password: Windows Systems

On Windows, use the following procedure to reset the password for the MySQL `'root'@'localhost'` account. To change the password for a `root` account with a different host name part, modify the instructions to use that host name.

1. Log on to your system as Administrator.
2. Stop the MySQL server if it is running. For a server that is running as a Windows service, go to the Services manager: From the **Start** menu, select **Control Panel**, then **Administrative Tools**, then **Services**. Find the MySQL service in the list and stop it.

   If your server is not running as a service, you may need to use the Task Manager to force it to stop.
3. Create a text file containing the password-assignment statement on a single line. Replace the password with the password that you want to use.

```
ALTER USER 'root'@'localhost' IDENTIFIED BY 'MyNewPass';
```

4. Save the file. This example assumes that you name the file `C:\mysql-init.txt`.

5. Open a console window to get to the command prompt: From the **Start** menu, select **Run**, then enter `cmd` as the command to be run.

6. Start the MySQL server with the special `--init-file` option (notice that the backslash in the option value is doubled):

   ```
   C:\> cd "C:\Program Files\MySQL\MySQL Server 8.0\bin"
   C:\> mysqld --init-file=C:\\mysql-init.txt
   ```

   If you installed MySQL to a different location, adjust the `cd` command accordingly.

   The server executes the contents of the file named by the `--init-file` option at startup, changing the `'root'@'localhost'` account password.

   To have server output to appear in the console window rather than in a log file, add the `--console` option to the `mysqld` command.

   If you installed MySQL using the MySQL Installation Wizard, you may need to specify a `--defaults-file` option. For example:

   ```
   C:\> mysqld
           --defaults-file="C:\\ProgramData\\MySQL\\MySQL Server 8.0\\my.ini"
           --init-file=C:\\mysql-init.txt
   ```

   The appropriate `--defaults-file` setting can be found using the Services Manager: From the **Start** menu, select **Control Panel**, then **Administrative Tools**, then **Services**. Find the MySQL service in the list, right-click it, and choose the `Properties` option. The `Path to executable` field contains the `--defaults-file` setting.

7. After the server has started successfully, delete `C:\mysql-init.txt`.

You should now be able to connect to the MySQL server as `root` using the new password. Stop the MySQL server and restart it normally. If you run the server as a service, start it from the Windows Services window. If you start the server manually, use whatever command you normally use.

Resetting the Root Password: Unix and Unix-Like Systems

On Unix, use the following procedure to reset the password for the MySQL `'root'@'localhost'` account. To change the password for a `root` account with a different host name part, modify the instructions to use that host name.

The instructions assume that you will start the MySQL server from the Unix login account that you normally use for running it. For example, if you run the server using the `mysql` login account, you should log in as `mysql` before using the instructions. Alternatively, you can log in as `root`, but in this case you *must* start `mysqld` with the `--user=mysql` option. If you start the server as `root` without using `--user=mysql`, the server may create `root`-owned files in the data directory, such as log files, and these may cause permission-related problems for future server startups. If that happens, you will need to either change the ownership of the files to `mysql` or remove them.

1. Log on to your system as the Unix user that the MySQL server runs as (for example, `mysql`).

2. Stop the MySQL server if it is running. Locate the `.pid` file that contains the server's process ID. The exact location and name of this file depend on your distribution, host name, and configuration. Common locations are `/var/lib/mysql/`, `/var/run/mysqld/`, and `/usr/local/mysql/data/`. Generally, the file name has an extension of `.pid` and begins with either `mysqld` or your system's host name.

Stop the MySQL server by sending a normal `kill` (not `kill -9`) to the `mysqld` process. Use the actual path name of the `.pid` file in the following command:

```
shell> kill `cat /mysql-data-directory/host_name.pid`
```

Use backticks (not forward quotation marks) with the `cat` command. These cause the output of `cat` to be substituted into the `kill` command.

3. Create a text file containing the password-assignment statement on a single line. Replace the password with the password that you want to use.

```
ALTER USER 'root'@'localhost' IDENTIFIED BY 'MyNewPass';
```

4. Save the file. This example assumes that you name the file `/home/me/mysql-init`. The file contains the password, so do not save it where it can be read by other users. If you are not logged in as `mysql` (the user the server runs as), make sure that the file has permissions that permit `mysql` to read it.

5. Start the MySQL server with the special `--init-file` option:

```
shell> mysqld --init-file=/home/me/mysql-init &
```

The server executes the contents of the file named by the `--init-file` option at startup, changing the `'root'@'localhost'` account password.

Other options may be necessary as well, depending on how you normally start your server. For example, `--defaults-file` may be needed before `--init-file`.

6. After the server has started successfully, delete `/home/me/mysql-init`.

You should now be able to connect to the MySQL server as `root` using the new password. Stop the server and restart it normally.

Resetting the Root Password: Generic Instructions

The preceding sections provide password-resetting instructions specifically for Windows and Unix and Unix-like systems. Alternatively, on any platform, you can reset the password using the `mysql` client (but this approach is less secure):

1. Stop the MySQL server if necessary, then restart it with the `--skip-grant-tables` option. This enables anyone to connect without a password and with all privileges, and disables account-management statements such as `ALTER USER` and `SET PASSWORD`. Because this is insecure, you might want to use `--skip-grant-tables` in conjunction with `--skip-networking` to prevent remote clients from connecting.

2. Connect to the MySQL server using the `mysql` client; no password is necessary because the server was started with `--skip-grant-tables`:

```
shell> mysql
```

3. In the `mysql` client, tell the server to reload the grant tables so that account-management statements work:

```
mysql> FLUSH PRIVILEGES;
```

Then change the `'root'@'localhost'` account password. Replace the password with the password that you want to use. To change the password for a `root` account with a different host name part, modify the instructions to use that host name.

```
mysql> ALTER USER 'root'@'localhost' IDENTIFIED BY 'MyNewPass';
```

You should now be able to connect to the MySQL server as `root` using the new password. Stop the server and restart it normally (without the `--skip-grant-tables` and `--skip-networking` options).

B.5.3.3 What to Do If MySQL Keeps Crashing

Each MySQL version is tested on many platforms before it is released. This does not mean that there are no bugs in MySQL, but if there are bugs, they should be very few and can be hard to find. If you have a problem, it always helps if you try to find out exactly what crashes your system, because you have a much better chance of getting the problem fixed quickly.

First, you should try to find out whether the problem is that the `mysqld` server dies or whether your problem has to do with your client. You can check how long your `mysqld` server has been up by executing `mysqladmin version`. If `mysqld` has died and restarted, you may find the reason by looking in the server's error log. See Section 5.4.2, “The Error Log”.

On some systems, you can find in the error log a stack trace of where `mysqld` died that you can resolve with the `resolve_stack_dump` program. See Section 28.5, “Debugging and Porting MySQL”. Note that the variable values written in the error log may not always be 100% correct.

Many server crashes are caused by corrupted data files or index files. MySQL updates the files on disk with the `write()` system call after every SQL statement and before the client is notified about the result. (This is not true if you are running with `--delay-key-write`, in which case data files are written but not index files.) This means that data file contents are safe even if `mysqld` crashes, because the operating system ensures that the unflushed data is written to disk. You can force MySQL to flush everything to disk after every SQL statement by starting `mysqld` with the `--flush` option.

The preceding means that normally you should not get corrupted tables unless one of the following happens:

- The MySQL server or the server host was killed in the middle of an update.
- You have found a bug in `mysqld` that caused it to die in the middle of an update.
- Some external program is manipulating data files or index files at the same time as `mysqld` without locking the table properly.
- You are running many `mysqld` servers using the same data directory on a system that does not support good file system locks (normally handled by the `lockd` lock manager), or you are running multiple servers with external locking disabled.
- You have a crashed data file or index file that contains very corrupt data that confused `mysqld`.
- You have found a bug in the data storage code. This isn't likely, but it is at least possible. In this case, you can try to change the storage engine to another engine by using `ALTER TABLE` on a repaired copy of the table.

Because it is very difficult to know why something is crashing, first try to check whether things that work for others crash for you. Try the following things:

- Stop the `mysqld` server with `mysqladmin shutdown`, run `myisamchk --silent --force */*.MYI` from the data directory to check all `MyISAM` tables, and restart `mysqld`. This ensures that you are running from a clean state. See Chapter 5, *MySQL Server Administration*.
- Start `mysqld` with the general query log enabled (see Section 5.4.3, “The General Query Log”). Then try to determine from the information written to the log whether some specific query kills the server. About 95% of all bugs are related to a particular query. Normally, this is one of the last

queries in the log file just before the server restarts. See Section 5.4.3, “The General Query Log”. If you can repeatedly kill MySQL with a specific query, even when you have checked all tables just before issuing it, then you have isolated the bug and should submit a bug report for it. See Section 1.7, “How to Report Bugs or Problems”.

- Try to make a test case that we can use to repeat the problem. See Section 28.5, “Debugging and Porting MySQL”.
- Try the `fork_big.pl` script. (It is located in the `tests` directory of source distributions.)
- Configuring MySQL for debugging makes it much easier to gather information about possible errors if something goes wrong. Reconfigure MySQL with the `-DWITH_DEBUG=1` option to `CMake` and then recompile. See Section 28.5, “Debugging and Porting MySQL”.
- Make sure that you have applied the latest patches for your operating system.
- Use the `--skip-external-locking` option to `mysqld`. On some systems, the `lockd` lock manager does not work properly; the `--skip-external-locking` option tells `mysqld` not to use external locking. (This means that you cannot run two `mysqld` servers on the same data directory and that you must be careful if you use `myisamchk`. Nevertheless, it may be instructive to try the option as a test.)
- If `mysqld` appears to be running but not responding, try `mysqladmin -u root processlist`. Sometimes `mysqld` is not hung even though it seems unresponsive. The problem may be that all connections are in use, or there may be some internal lock problem. `mysqladmin -u root processlist` usually is able to make a connection even in these cases, and can provide useful information about the current number of connections and their status.
- Run the command `mysqladmin -i 5 status` or `mysqladmin -i 5 -r status` in a separate window to produce statistics while running other queries.
- Try the following:

  1. Start `mysqld` from `gdb` (or another debugger). See Section 28.5, “Debugging and Porting MySQL”.
  2. Run your test scripts.
  3. Print the backtrace and the local variables at the three lowest levels. In `gdb`, you can do this with the following commands when `mysqld` has crashed inside `gdb`:

     ```
     backtrace
     info local
     up
     info local
     up
     info local
     ```

     With `gdb`, you can also examine which threads exist with `info threads` and switch to a specific thread with `thread N`, where `N` is the thread ID.

- Try to simulate your application with a Perl script to force MySQL to crash or misbehave.
- Send a normal bug report. See Section 1.7, “How to Report Bugs or Problems”. Be even more detailed than usual. Because MySQL works for many people, the crash might result from something that exists only on your computer (for example, an error that is related to your particular system libraries).
- If you have a problem with tables containing dynamic-length rows and you are using only `VARCHAR` columns (not `BLOB` or `TEXT` columns), you can try to change all `VARCHAR` to `CHAR` with `ALTER TABLE`. This forces MySQL to use fixed-size rows. Fixed-size rows take a little extra space, but are much more tolerant to corruption.

The current dynamic row code has been in use for several years with very few problems, but dynamic-length rows are by nature more prone to errors, so it may be a good idea to try this strategy to see whether it helps.

- Consider the possibility of hardware faults when diagnosing problems. Defective hardware can be the cause of data corruption. Pay particular attention to your memory and disk subsystems when troubleshooting hardware.

B.5.3.4 How MySQL Handles a Full Disk

This section describes how MySQL responds to disk-full errors (such as “no space left on device”), and to quota-exceeded errors (such as “write failed” or “user block limit reached”).

This section is relevant for writes to `MyISAM` tables. It also applies for writes to binary log files and binary log index file, except that references to “row” and “record” should be understood to mean “event.”

When a disk-full condition occurs, MySQL does the following:

- It checks once every minute to see whether there is enough space to write the current row. If there is enough space, it continues as if nothing had happened.
- Every 10 minutes it writes an entry to the log file, warning about the disk-full condition.

To alleviate the problem, take the following actions:

- To continue, you only have to free enough disk space to insert all records.
- Alternatively, to abort the thread, use `mysqladmin kill`. The thread is aborted the next time it checks the disk (in one minute).
- Other threads might be waiting for the table that caused the disk-full condition. If you have several “locked” threads, killing the one thread that is waiting on the disk-full condition enables the other threads to continue.

Exceptions to the preceding behavior are when you use `REPAIR TABLE` or `OPTIMIZE TABLE` or when the indexes are created in a batch after `LOAD DATA INFILE` or after an `ALTER TABLE` statement. All of these statements may create large temporary files that, if left to themselves, would cause big problems for the rest of the system. If the disk becomes full while MySQL is doing any of these operations, it removes the big temporary files and mark the table as crashed. The exception is that for `ALTER TABLE`, the old table is left unchanged.

B.5.3.5 Where MySQL Stores Temporary Files

On Unix, MySQL uses the value of the `TMPDIR` environment variable as the path name of the directory in which to store temporary files. If `TMPDIR` is not set, MySQL uses the system default, which is usually `/tmp`, `/var/tmp`, or `/usr/tmp`.

On Windows, MySQL checks in order the values of the `TMPDIR`, `TEMP`, and `TMP` environment variables. For the first one found to be set, MySQL uses it and does not check those remaining. If none of `TMPDIR`, `TEMP`, or `TMP` are set, MySQL uses the Windows system default, which is usually `C:\windows\temp\`.

If the file system containing your temporary file directory is too small, you can use the `mysqld --tmpdir` option to specify a directory in a file system where you have enough space. On replication slaves, you can use `--slave-load-tmpdir` to specify a separate directory for holding temporary files when replicating `LOAD DATA INFILE` statements.

The `--tmpdir` option can be set to a list of several paths that are used in round-robin fashion. Paths should be separated by colon characters (`:`) on Unix and semicolon characters (`;`) on Windows.

Note

To spread the load effectively, these paths should be located on different *physical* disks, not different partitions of the same disk.

If the MySQL server is acting as a replication slave, you should be sure to set `--slave-load-tmpdir` not to point to a directory that is on a memory-based file system or to a directory that is cleared when the server host restarts. A replication slave needs some of its temporary files to survive a machine restart so that it can replicate temporary tables or `LOAD DATA INFILE` operations. If files in the slave temporary file directory are lost when the server restarts, replication fails.

MySQL arranges that temporary files are removed if `mysqld` is terminated. On platforms that support it (such as Unix), this is done by unlinking the file after opening it. The disadvantage of this is that the name does not appear in directory listings and you do not see a big temporary file that fills up the file system in which the temporary file directory is located. (In such cases, `lsof +L1` may be helpful in identifying large files associated with `mysqld`.)

When sorting (`ORDER BY` or `GROUP BY`), MySQL normally uses one or two temporary files. The maximum disk space required is determined by the following expression:

```
(length of what is sorted + sizeof(row pointer))
* number of matched rows
* 2
```

The row pointer size is usually four bytes, but may grow in the future for really big tables.

For some `SELECT` queries, MySQL also creates temporary SQL tables. These are not hidden and have names of the form `SQL_*`.

In most cases, `ALTER TABLE` creates a temporary copy of the original table in the same directory as the original table.

Where InnoDB Stores Temporary Files

`InnoDB` temporary tables are stored in a temporary tablespace file named `ibtmp1`, which is located in the MySQL `data` directory (`datadir`) by default. The `innodb_temp_data_file_path` option can be used at startup to specify a different file name and location.

If an `ALTER TABLE` operation on an `InnoDB` table uses the `ALGORITHM=COPY` technique, `InnoDB` creates a temporary copy of the table in the same directory as the original table. Temporary table file names begin with an `#sql-` prefix and only appear briefly during the `ALTER TABLE` operation.

If an `ALTER TABLE` operation rebuilds an `InnoDB` table using the `ALGORITHM=INPLACE` technique (online DDL), `InnoDB` creates an *intermediate* copy of the table the same directory as the original table. Intermediate table file names begin with an `#sql-ib` prefix and only appear briefly during the `ALTER TABLE` operation.

`ALTER TABLE` operations that rebuild an `InnoDB` table using the `ALGORITHM=INPLACE` technique (online DDL) also create temporary sort files in the MySQL temporary directory (`$TMPDIR` on Unix, `%TEMP%` on Windows, or the directory specified by the `--tmpdir` configuration option). If the temporary directory is not large enough to hold such files, you may need to reconfigure `tmpdir`. Alternatively, you can define a separate temporary directory for `InnoDB` online `ALTER TABLE` operations using the `innodb_tmpdir` configuration option. This option was introduced to help avoid temporary directory overflows that could occur as a result of large temporary sort files created during online `ALTER TABLE` operations. `innodb_tmpdir` can be configured dynamically using a `SET GLOBAL` or `SET SESSION` statement.

The `innodb_tmpdir` option is not applicable to intermediate table files, which are always created in the same directory as the original table.

In replication environments, only consider replicating an `innodb_tmpdir` setting if all servers have the same operating system environment. Otherwise, replicating an `innodb_tmpdir` setting could result in a replication failure when running online `ALTER TABLE` operations. If server operating environments differ, it is recommended that you configure `innodb_tmpdir` on each server individually.

For more information about online DDL, Section 15.12, “InnoDB and Online DDL”.

B.5.3.6 How to Protect or Change the MySQL Unix Socket File

The default location for the Unix socket file that the server uses for communication with local clients is `/tmp/mysql.sock`. (For some distribution formats, the directory might be different, such as `/var/lib/mysql` for RPMs.)

On some versions of Unix, anyone can delete files in the `/tmp` directory or other similar directories used for temporary files. If the socket file is located in such a directory on your system, this might cause problems.

On most versions of Unix, you can protect your `/tmp` directory so that files can be deleted only by their owners or the superuser (`root`). To do this, set the `sticky` bit on the `/tmp` directory by logging in as `root` and using the following command:

```
shell> chmod +t /tmp
```

You can check whether the `sticky` bit is set by executing `ls -ld /tmp`. If the last permission character is `t`, the bit is set.

Another approach is to change the place where the server creates the Unix socket file. If you do this, you should also let client programs know the new location of the file. You can specify the file location in several ways:

- Specify the path in a global or local option file. For example, put the following lines in `/etc/my.cnf`:

  ```
  [mysqld]
  socket=/path/to/socket

  [client]
  socket=/path/to/socket
  ```

  See Section 4.2.6, “Using Option Files”.

- Specify a `--socket` option on the command line to `mysqld_safe` and when you run client programs.

- Set the `MYSQL_UNIX_PORT` environment variable to the path of the Unix socket file.

- Recompile MySQL from source to use a different default Unix socket file location. Define the path to the file with the `MYSQL_UNIX_ADDR` option when you run `CMake`. See Section 2.8.4, “MySQL Source-Configuration Options”.

You can test whether the new socket location works by attempting to connect to the server with this command:

```
shell> mysqladmin --socket=/path/to/socket version
```

B.5.3.7 Time Zone Problems

If you have a problem with `SELECT NOW()` returning values in UTC and not your local time, you have to tell the server your current time zone. The same applies if `UNIX_TIMESTAMP()` returns the wrong value. This should be done for the environment in which the server runs; for example, in `mysqld_safe` or `mysql.server`. See Section 4.9, “MySQL Program Environment Variables”.

You can set the time zone for the server with the `--timezone=timezone_name` option to `mysqld_safe`. You can also set it by setting the `TZ` environment variable before you start `mysqld`.

The permissible values for `--timezone` or `TZ` are system dependent. Consult your operating system documentation to see what values are acceptable.

B.5.4 Query-Related Issues

B.5.4.1 Case Sensitivity in String Searches

For nonbinary strings (`CHAR`, `VARCHAR`, `TEXT`), string searches use the collation of the comparison operands. For binary strings (`BINARY`, `VARBINARY`, `BLOB`), comparisons use the numeric values of the bytes in the operands; this means that for alphabetic characters, comparisons will be case sensitive.

A comparison between a nonbinary string and binary string is treated as a comparison of binary strings.

Simple comparison operations (`>=, >, =, <, <=`, sorting, and grouping) are based on each character's "sort value." Characters with the same sort value are treated as the same character. For example, if `e` and `é` have the same sort value in a given collation, they compare as equal.

The default character set and collation are `utf8mb4` and `utf8mb4_0900_ai_ci`, so nonbinary string comparisons are case insensitive by default. This means that if you search with *`col_name`* `LIKE 'a%'`, you get all column values that start with `A` or `a`. To make this search case sensitive, make sure that one of the operands has a case sensitive or binary collation. For example, if you are comparing a column and a string that both have the `utf8mb4` character set, you can use the `COLLATE` operator to cause either operand to have the `utf8mb4_0900_as_cs` or `utf8mb4_bin` collation:

```
col_name COLLATE utf8mb4_0900_as_cs LIKE 'a%'
col_name LIKE 'a%' COLLATE utf8mb4_0900_as_cs
col_name COLLATE utf8mb4_bin LIKE 'a%'
col_name LIKE 'a%' COLLATE utf8mb4_bin
```

If you want a column always to be treated in case-sensitive fashion, declare it with a case sensitive or binary collation. See Section 13.1.15, "CREATE TABLE Syntax".

To cause a case-sensitive comparison of nonbinary strings to be case insensitive, use `COLLATE` to name a case-insensitive collation. The strings in the following example normally are case sensitive, but `COLLATE` changes the comparison to be case insensitive:

```
mysql> SET NAMES 'utf8mb4';
mysql> SET @s1 = 'MySQL' COLLATE utf8mb4_bin,
           @s2 = 'mysql' COLLATE utf8mb4_bin;
mysql> SELECT @s1 = @s2;
+-----------+
| @s1 = @s2 |
+-----------+
|         0 |
+-----------+
mysql> SELECT @s1 COLLATE utf8mb4_0900_ai_ci = @s2;
+--------------------------------------+
| @s1 COLLATE utf8mb4_0900_ai_ci = @s2 |
+--------------------------------------+
|                                    1 |
+--------------------------------------+
```

A binary string is case sensitive in comparisons. To compare the string as case insensitive, convert it to a nonbinary string and use `COLLATE` to name a case-insensitive collation:

```
mysql> SET @s = BINARY 'MySQL';
mysql> SELECT @s = 'mysql';
+--------------+
| @s = 'mysql' |
```

```
+--------------+
|            0 |
+--------------+
mysql> SELECT CONVERT(@s USING utf8mb4) COLLATE utf8mb4_0900_ai_ci = 'mysql';
+-------------------------------------------------------------------+
| CONVERT(@s USING utf8mb4) COLLATE utf8mb4_0900_ai_ci = 'mysql' |
+-------------------------------------------------------------------+
|                                                                 1 |
+-------------------------------------------------------------------+
```

To determine whether a value will compare as a nonbinary or binary string, use the `COLLATION()` function. This example shows that `VERSION()` returns a string that has a case-insensitive collation, so comparisons are case insensitive:

```
mysql> SELECT COLLATION(VERSION());
+----------------------+
| COLLATION(VERSION()) |
+----------------------+
| utf8_general_ci      |
+----------------------+
```

For binary strings, the collation value is `binary`, so comparisons will be case sensitive. One context in which you will see `binary` is for compression and encryption functions, which return binary strings as a general rule: string:

```
mysql> SELECT COLLATION(ENCRYPT('x')), COLLATION(SHA1('x'));
+-------------------------+----------------------+
| COLLATION(ENCRYPT('x')) | COLLATION(SHA1('x')) |
+-------------------------+----------------------+
| binary                  | binary               |
+-------------------------+----------------------+
```

To check the sort value of a string, the `WEIGHT_STRING()` may be helpful. See Section 12.5, “String Functions”.

B.5.4.2 Problems Using DATE Columns

The format of a `DATE` value is `'YYYY-MM-DD'`. According to standard SQL, no other format is permitted. You should use this format in `UPDATE` expressions and in the `WHERE` clause of `SELECT` statements. For example:

```
SELECT * FROM t1 WHERE date >= '2003-05-05';
```

As a convenience, MySQL automatically converts a date to a number if the date is used in a numeric context and vice versa. MySQL also permits a “relaxed” string format when updating and in a `WHERE` clause that compares a date to a `DATE`, `DATETIME`, or `TIMESTAMP` column. “Relaxed” format means that any punctuation character may be used as the separator between parts. For example, `'2004-08-15'` and `'2004#08#15'` are equivalent. MySQL can also convert a string containing no separators (such as `'20040815'`), provided it makes sense as a date.

When you compare a `DATE`, `TIME`, `DATETIME`, or `TIMESTAMP` to a constant string with the `<`, `<=`, `=`, `>=`, `>`, or `BETWEEN` operators, MySQL normally converts the string to an internal long integer for faster comparison (and also for a bit more “relaxed” string checking). However, this conversion is subject to the following exceptions:

- When you compare two columns
- When you compare a `DATE`, `TIME`, `DATETIME`, or `TIMESTAMP` column to an expression
- When you use any comparison method other than those just listed, such as `IN` or `STRCMP()`.

For those exceptions, the comparison is done by converting the objects to strings and performing a string comparison.

To be on the safe side, assume that strings are compared as strings and use the appropriate string functions if you want to compare a temporal value to a string.

The special “zero” date `'0000-00-00'` can be stored and retrieved as `'0000-00-00'`. When a `'0000-00-00'` date is used through Connector/ODBC, it is automatically converted to `NULL` because ODBC cannot handle that kind of date.

Because MySQL performs the conversions just described, the following statements work (assume that `idate` is a `DATE` column):

```
INSERT INTO t1 (idate) VALUES (19970505);
INSERT INTO t1 (idate) VALUES ('19970505');
INSERT INTO t1 (idate) VALUES ('97-05-05');
INSERT INTO t1 (idate) VALUES ('1997.05.05');
INSERT INTO t1 (idate) VALUES ('1997 05 05');
INSERT INTO t1 (idate) VALUES ('0000-00-00');

SELECT idate FROM t1 WHERE idate >= '1997-05-05';
SELECT idate FROM t1 WHERE idate >= 19970505;
SELECT MOD(idate,100) FROM t1 WHERE idate >= 19970505;
SELECT idate FROM t1 WHERE idate >= '19970505';
```

However, the following statement does not work:

```
SELECT idate FROM t1 WHERE STRCMP(idate,'20030505')=0;
```

`STRCMP()` is a string function, so it converts `idate` to a string in `'YYYY-MM-DD'` format and performs a string comparison. It does not convert `'20030505'` to the date `'2003-05-05'` and perform a date comparison.

If you enable the `ALLOW_INVALID_DATES` SQL mode, MySQL permits you to store dates that are given only limited checking: MySQL requires only that the day is in the range from 1 to 31 and the month is in the range from 1 to 12. This makes MySQL very convenient for Web applications where you obtain year, month, and day in three different fields and you want to store exactly what the user inserted (without date validation).

MySQL permits you to store dates where the day or month and day are zero. This is convenient if you want to store a birthdate in a `DATE` column and you know only part of the date. To disallow zero month or day parts in dates, enable the `NO_ZERO_IN_DATE` mode.

MySQL permits you to store a “zero” value of `'0000-00-00'` as a “dummy date.” This is in some cases more convenient than using `NULL` values. If a date to be stored in a `DATE` column cannot be converted to any reasonable value, MySQL stores `'0000-00-00'`. To disallow `'0000-00-00'`, enable the `NO_ZERO_DATE` mode.

To have MySQL check all dates and accept only legal dates (unless overridden by `IGNORE`), set the `sql_mode` system variable to `"NO_ZERO_IN_DATE,NO_ZERO_DATE"`.

B.5.4.3 Problems with NULL Values

The concept of the `NULL` value is a common source of confusion for newcomers to SQL, who often think that `NULL` is the same thing as an empty string `''`. This is not the case. For example, the following statements are completely different:

```
mysql> INSERT INTO my_table (phone) VALUES (NULL);
mysql> INSERT INTO my_table (phone) VALUES ('');
```

Both statements insert a value into the `phone` column, but the first inserts a `NULL` value and the second inserts an empty string. The meaning of the first can be regarded as “phone number is not known” and the meaning of the second can be regarded as “the person is known to have no phone, and thus no phone number.”

To help with NULL handling, you can use the IS NULL and IS NOT NULL operators and the IFNULL() function.

In SQL, the NULL value is never true in comparison to any other value, even NULL. An expression that contains NULL always produces a NULL value unless otherwise indicated in the documentation for the operators and functions involved in the expression. All columns in the following example return NULL:

```
mysql> SELECT NULL, 1+NULL, CONCAT('Invisible',NULL);
```

To search for column values that are NULL, you cannot use an expr = NULL test. The following statement returns no rows, because expr = NULL is never true for any expression:

```
mysql> SELECT * FROM my_table WHERE phone = NULL;
```

To look for NULL values, you must use the IS NULL test. The following statements show how to find the NULL phone number and the empty phone number:

```
mysql> SELECT * FROM my_table WHERE phone IS NULL;
mysql> SELECT * FROM my_table WHERE phone = '';
```

See Section 3.3.4.6, “Working with NULL Values”, for additional information and examples.

You can add an index on a column that can have NULL values if you are using the MyISAM, InnoDB, or MEMORY storage engine. Otherwise, you must declare an indexed column NOT NULL, and you cannot insert NULL into the column.

When reading data with LOAD DATA INFILE, empty or missing columns are updated with ''. To load a NULL value into a column, use \N in the data file. The literal word NULL may also be used under some circumstances. See Section 13.2.7, “LOAD DATA INFILE Syntax”.

When using DISTINCT, GROUP BY, or ORDER BY, all NULL values are regarded as equal.

When using ORDER BY, NULL values are presented first, or last if you specify DESC to sort in descending order.

Aggregate (summary) functions such as COUNT(), MIN(), and SUM() ignore NULL values. The exception to this is COUNT(*), which counts rows and not individual column values. For example, the following statement produces two counts. The first is a count of the number of rows in the table, and the second is a count of the number of non-NULL values in the age column:

```
mysql> SELECT COUNT(*), COUNT(age) FROM person;
```

For some data types, MySQL handles NULL values specially. If you insert NULL into a TIMESTAMP column, the current date and time is inserted. If you insert NULL into an integer or floating-point column that has the AUTO_INCREMENT attribute, the next number in the sequence is inserted.

B.5.4.4 Problems with Column Aliases

An alias can be used in a query select list to give a column a different name. You can use the alias in GROUP BY, ORDER BY, or HAVING clauses to refer to the column:

```
SELECT SQRT(a*b) AS root FROM tbl_name
  GROUP BY root HAVING root > 0;
SELECT id, COUNT(*) AS cnt FROM tbl_name
  GROUP BY id HAVING cnt > 0;
SELECT id AS 'Customer identity' FROM tbl_name;
```

Standard SQL disallows references to column aliases in a WHERE clause. This restriction is imposed because when the WHERE clause is evaluated, the column value may not yet have been determined. For example, the following query is illegal:

```
SELECT id, COUNT(*) AS cnt FROM tbl_name
  WHERE cnt > 0 GROUP BY id;
```

The `WHERE` clause determines which rows should be included in the `GROUP BY` clause, but it refers to the alias of a column value that is not known until after the rows have been selected, and grouped by the `GROUP BY`.

In the select list of a query, a quoted column alias can be specified using identifier or string quoting characters:

```
SELECT 1 AS `one`, 2 AS 'two';
```

Elsewhere in the statement, quoted references to the alias must use identifier quoting or the reference is treated as a string literal. For example, this statement groups by the values in column `id`, referenced using the alias `` `a` ``:

```
SELECT id AS 'a', COUNT(*) AS cnt FROM tbl_name
  GROUP BY `a`;
```

But this statement groups by the literal string `'a'` and will not work as expected:

```
SELECT id AS 'a', COUNT(*) AS cnt FROM tbl_name
  GROUP BY 'a';
```

B.5.4.5 Rollback Failure for Nontransactional Tables

If you receive the following message when trying to perform a `ROLLBACK`, it means that one or more of the tables you used in the transaction do not support transactions:

```
Warning: Some non-transactional changed tables couldn't be rolled back
```

These nontransactional tables are not affected by the `ROLLBACK` statement.

If you were not deliberately mixing transactional and nontransactional tables within the transaction, the most likely cause for this message is that a table you thought was transactional actually is not. This can happen if you try to create a table using a transactional storage engine that is not supported by your `mysqld` server (or that was disabled with a startup option). If `mysqld` does not support a storage engine, it instead creates the table as a `MyISAM` table, which is nontransactional.

You can check the storage engine for a table by using either of these statements:

```
SHOW TABLE STATUS LIKE 'tbl_name';
SHOW CREATE TABLE tbl_name;
```

See Section 13.7.5.36, “SHOW TABLE STATUS Syntax”, and Section 13.7.5.10, “SHOW CREATE TABLE Syntax”.

To check which storage engines your `mysqld` server supports, use this statement:

```
SHOW ENGINES;
```

See Section 13.7.5.16, “SHOW ENGINES Syntax” for full details.

B.5.4.6 Deleting Rows from Related Tables

If the total length of the `DELETE` statement for `related_table` is more than 1MB (the default value of the `max_allowed_packet` system variable), you should split it into smaller parts and execute

multiple `DELETE` statements. You probably get the fastest `DELETE` by specifying only 100 to 1,000 `related_column` values per statement if the `related_column` is indexed. If the `related_column` isn't indexed, the speed is independent of the number of arguments in the `IN` clause.

B.5.4.7 Solving Problems with No Matching Rows

If you have a complicated query that uses many tables but that returns no rows, you should use the following procedure to find out what is wrong:

1. Test the query with `EXPLAIN` to check whether you can find something that is obviously wrong. See Section 13.8.2, “EXPLAIN Syntax”.

2. Select only those columns that are used in the `WHERE` clause.

3. Remove one table at a time from the query until it returns some rows. If the tables are large, it is a good idea to use `LIMIT 10` with the query.

4. Issue a `SELECT` for the column that should have matched a row against the table that was last removed from the query.

5. If you are comparing `FLOAT` or `DOUBLE` columns with numbers that have decimals, you cannot use equality (`=`) comparisons. This problem is common in most computer languages because not all floating-point values can be stored with exact precision. In some cases, changing the `FLOAT` to a `DOUBLE` fixes this. See Section B.5.4.8, “Problems with Floating-Point Values”.

6. If you still cannot figure out what is wrong, create a minimal test that can be run with `mysql test < query.sql` that shows your problems. You can create a test file by dumping the tables with `mysqldump --quick db_name tbl_name_1 ... tbl_name_n > query.sql`. Open the file in an editor, remove some insert lines (if there are more than needed to demonstrate the problem), and add your `SELECT` statement at the end of the file.

   Verify that the test file demonstrates the problem by executing these commands:

   ```
   shell> mysqladmin create test2
   shell> mysql test2 < query.sql
   ```

   Attach the test file to a bug report, which you can file using the instructions in Section 1.7, “How to Report Bugs or Problems”.

B.5.4.8 Problems with Floating-Point Values

Floating-point numbers sometimes cause confusion because they are approximate and not stored as exact values. A floating-point value as written in an SQL statement may not be the same as the value represented internally. Attempts to treat floating-point values as exact in comparisons may lead to problems. They are also subject to platform or implementation dependencies. The `FLOAT` and `DOUBLE` data types are subject to these issues. For `DECIMAL` columns, MySQL performs operations with a precision of 65 decimal digits, which should solve most common inaccuracy problems.

The following example uses `DOUBLE` to demonstrate how calculations that are done using floating-point operations are subject to floating-point error.

```
mysql> CREATE TABLE t1 (i INT, d1 DOUBLE, d2 DOUBLE);
mysql> INSERT INTO t1 VALUES (1, 101.40, 21.40), (1, -80.00, 0.00),
    -> (2, 0.00, 0.00), (2, -13.20, 0.00), (2, 59.60, 46.40),
    -> (2, 30.40, 30.40), (3, 37.00, 7.40), (3, -29.60, 0.00),
    -> (4, 60.00, 15.40), (4, -10.60, 0.00), (4, -34.00, 0.00),
    -> (5, 33.00, 0.00), (5, -25.80, 0.00), (5, 0.00, 7.20),
    -> (6, 0.00, 0.00), (6, -51.40, 0.00);

mysql> SELECT i, SUM(d1) AS a, SUM(d2) AS b
    -> FROM t1 GROUP BY i HAVING a <> b;
```

```
+------+-------+------+
| i    | a     | b    |
+------+-------+------+
|    1 |  21.4 | 21.4 |
|    2 |  76.8 | 76.8 |
|    3 |   7.4 |  7.4 |
|    4 |  15.4 | 15.4 |
|    5 |   7.2 |  7.2 |
|    6 | -51.4 |    0 |
+------+-------+------+
```

The result is correct. Although the first five records look like they should not satisfy the comparison (the values of `a` and `b` do not appear to be different), they may do so because the difference between the numbers shows up around the tenth decimal or so, depending on factors such as computer architecture or the compiler version or optimization level. For example, different CPUs may evaluate floating-point numbers differently.

If columns `d1` and `d2` had been defined as `DECIMAL` rather than `DOUBLE`, the result of the `SELECT` query would have contained only one row—the last one shown above.

The correct way to do floating-point number comparison is to first decide on an acceptable tolerance for differences between the numbers and then do the comparison against the tolerance value. For example, if we agree that floating-point numbers should be regarded the same if they are same within a precision of one in ten thousand (0.0001), the comparison should be written to find differences larger than the tolerance value:

```
mysql> SELECT i, SUM(d1) AS a, SUM(d2) AS b FROM t1
    -> GROUP BY i HAVING ABS(a - b) > 0.0001;
+------+-------+------+
| i    | a     | b    |
+------+-------+------+
|    6 | -51.4 |    0 |
+------+-------+------+
1 row in set (0.00 sec)
```

Conversely, to get rows where the numbers are the same, the test should find differences within the tolerance value:

```
mysql> SELECT i, SUM(d1) AS a, SUM(d2) AS b FROM t1
    -> GROUP BY i HAVING ABS(a - b) <= 0.0001;
+------+------+------+
| i    | a    | b    |
+------+------+------+
|    1 | 21.4 | 21.4 |
|    2 | 76.8 | 76.8 |
|    3 |  7.4 |  7.4 |
|    4 | 15.4 | 15.4 |
|    5 |  7.2 |  7.2 |
+------+------+------+
5 rows in set (0.03 sec)
```

Floating-point values are subject to platform or implementation dependencies. Suppose that you execute the following statements:

```
CREATE TABLE t1(c1 FLOAT(53,0), c2 FLOAT(53,0));
INSERT INTO t1 VALUES('1e+52','-1e+52');
SELECT * FROM t1;
```

On some platforms, the `SELECT` statement returns `inf` and `-inf`. On others, it returns `0` and `-0`.

An implication of the preceding issues is that if you attempt to create a replication slave by dumping table contents with `mysqldump` on the master and reloading the dump file into the slave, tables containing floating-point columns might differ between the two hosts.

B.5.5 Optimizer-Related Issues

MySQL uses a cost-based optimizer to determine the best way to resolve a query. In many cases, MySQL can calculate the best possible query plan, but sometimes MySQL does not have enough information about the data at hand and has to make “educated” guesses about the data.

For the cases when MySQL does not do the "right" thing, tools that you have available to help MySQL are:

- Use the `EXPLAIN` statement to get information about how MySQL processes a query. To use it, just add the keyword `EXPLAIN` to the front of your `SELECT` statement:

  ```
  mysql> EXPLAIN SELECT * FROM t1, t2 WHERE t1.i = t2.i;
  ```

  `EXPLAIN` is discussed in more detail in Section 13.8.2, “EXPLAIN Syntax”.

- Use `ANALYZE TABLE tbl_name` to update the key distributions for the scanned table. See Section 13.7.2.1, “ANALYZE TABLE Syntax”.

- Use `FORCE INDEX` for the scanned table to tell MySQL that table scans are very expensive compared to using the given index:

  ```
  SELECT * FROM t1, t2 FORCE INDEX (index_for_column)
  WHERE t1.col_name=t2.col_name;
  ```

  `USE INDEX` and `IGNORE INDEX` may also be useful. See Section 8.9.4, “Index Hints”.

- Global and table-level `STRAIGHT_JOIN`. See Section 13.2.10, “SELECT Syntax”.

- You can tune global or thread-specific system variables. For example, start `mysqld` with the `--max-seeks-for-key=1000` option or use `SET max_seeks_for_key=1000` to tell the optimizer to assume that no key scan causes more than 1,000 key seeks. See Section 5.1.5, “Server System Variables”.

B.5.6 Table Definition-Related Issues

B.5.6.1 Problems with ALTER TABLE

If you get a duplicate-key error when using `ALTER TABLE` to change the character set or collation of a character column, the cause is either that the new column collation maps two keys to the same value or that the table is corrupted. In the latter case, you should run `REPAIR TABLE` on the table. `REPAIR TABLE` works for `MyISAM`, `ARCHIVE`, and `CSV` tables.

If you use `ALTER TABLE` on a transactional table or if you are using Windows, `ALTER TABLE` unlocks the table if you had done a `LOCK TABLE` on it. This is done because `InnoDB` and these operating systems cannot drop a table that is in use.

B.5.6.2 TEMPORARY Table Problems

Temporary tables created with `CREATE TEMPORARY TABLE` have the following limitations:

- `TEMPORARY` tables are supported only by the `InnoDB`, `MEMORY`, `MyISAM`, and `MERGE` storage engines.

- The `SHOW TABLES` statement does not list `TEMPORARY` tables.

- You cannot use `RENAME TABLE` to rename a `TEMPORARY` table. However, you can use `ALTER TABLE` instead:

  ```
  ALTER TABLE old_name RENAME new_name;
  ```

- You cannot refer to a `TEMPORARY` table more than once in the same query. For example, the following does not work:

```
SELECT * FROM temp_table JOIN temp_table AS t2;
```

  The statement produces this error:

```
ERROR 1137: Can't reopen table: 'temp_table'
```

  You can work around this issue if your query permits use of a common table expression (CTE) rather than a `TEMPORARY` table. For example, this fails with the `Can't reopen table` error:

```
CREATE TEMPORARY TABLE t SELECT 1 AS col_a, 2 AS col_b;
SELECT * FROM t AS t1 JOIN t AS t2;
```

  To avoid the error, use a `WITH` clause that defines a CTE, rather than the `TEMPORARY` table:

```
WITH cte AS (SELECT 1 AS col_a, 2 AS col_b)
SELECT * FROM cte AS t1 JOIN cte AS t2;
```

- The `Can't reopen table` error also occurs if you refer to a temporary table multiple times in a stored function under different aliases, even if the references occur in different statements within the function. It may occur for temporary tables created outside stored functions and referred to across multiple calling and callee functions.
- There are known issues in using temporary tables with replication. See Section 18.4.1.24, “Replication and Temporary Tables”, for more information.

B.5.7 Known Issues in MySQL

This section lists known issues in recent versions of MySQL.

For information about platform-specific issues, see the installation and porting instructions in Section 2.1, “General Installation Guidance”, and Section 28.5, “Debugging and Porting MySQL”.

The following problems are known:

- Subquery optimization for `IN` is not as effective as for `=`.
- Even if you use `lower_case_table_names=2` (which enables MySQL to remember the case used for databases and table names), MySQL does not remember the case used for database names for the function `DATABASE()` or within the various logs (on case-insensitive systems).
- Dropping a `FOREIGN KEY` constraint does not work in replication because the constraint may have another name on the slave.
- `REPLACE` (and `LOAD DATA` with the `REPLACE` option) does not trigger `ON DELETE CASCADE`.
- `DISTINCT` with `ORDER BY` does not work inside `GROUP_CONCAT()` if you do not use all and only those columns that are in the `DISTINCT` list.
- When inserting a big integer value (between 2^{63} and $2^{64}-1$) into a decimal or string column, it is inserted as a negative value because the number is evaluated in a signed integer context.
- With statement-based binary logging, the master writes the executed queries to the binary log. This is a very fast, compact, and efficient logging method that works perfectly in most cases. However, it is possible for the data on the master and slave to become different if a query is designed in such a way that the data modification is nondeterministic (generally not a recommended practice, even outside of replication).

For example:

- `CREATE TABLE ... SELECT` or `INSERT ... SELECT` statements that insert zero or `NULL` values into an `AUTO_INCREMENT` column.
- `DELETE` if you are deleting rows from a table that has foreign keys with `ON DELETE CASCADE` properties.
- `REPLACE ... SELECT`, `INSERT IGNORE ... SELECT` if you have duplicate key values in the inserted data.

If and only if the preceding queries have no `ORDER BY` clause guaranteeing a deterministic order.

For example, for `INSERT ... SELECT` with no `ORDER BY`, the `SELECT` may return rows in a different order (which results in a row having different ranks, hence getting a different number in the `AUTO_INCREMENT` column), depending on the choices made by the optimizers on the master and slave.

A query is optimized differently on the master and slave only if:

- The table is stored using a different storage engine on the master than on the slave. (It is possible to use different storage engines on the master and slave. For example, you can use `InnoDB` on the master, but `MyISAM` on the slave if the slave has less available disk space.)
- MySQL buffer sizes (`key_buffer_size`, and so on) are different on the master and slave.
- The master and slave run different MySQL versions, and the optimizer code differs between these versions.

This problem may also affect database restoration using `mysqlbinlog|mysql`.

The easiest way to avoid this problem is to add an `ORDER BY` clause to the aforementioned nondeterministic queries to ensure that the rows are always stored or modified in the same order. Using row-based or mixed logging format also avoids the problem.

- Log file names are based on the server host name if you do not specify a file name with the startup option. To retain the same log file names if you change your host name to something else, you must explicitly use options such as `--log-bin=`*`old_host_name`*`-bin`. See Section 5.1.4, “Server Command Options”. Alternatively, rename the old files to reflect your host name change. If these are binary logs, you must edit the binary log index file and fix the binary log file names there as well. (The same is true for the relay logs on a slave server.)
- `mysqlbinlog` does not delete temporary files left after a `LOAD DATA INFILE` statement. See Section 4.6.8, “`mysqlbinlog` — Utility for Processing Binary Log Files”.
- `RENAME` does not work with `TEMPORARY` tables or tables used in a `MERGE` table.
- When using `SET CHARACTER SET`, you cannot use translated characters in database, table, and column names.
- You cannot use `_` or `%` with `ESCAPE` in `LIKE ... ESCAPE`.
- The server uses only the first `max_sort_length` bytes when comparing data values. This means that values cannot reliably be used in `GROUP BY`, `ORDER BY`, or `DISTINCT` if they differ only after the first `max_sort_length` bytes. To work around this, increase the variable value. The default value of `max_sort_length` is 1024 and can be changed at server startup time or at runtime.
- Numeric calculations are done with `BIGINT` or `DOUBLE` (both are normally 64 bits long). Which precision you get depends on the function. The general rule is that bit functions are performed with `BIGINT` precision, `IF()` and `ELT()` with `BIGINT` or `DOUBLE` precision, and the rest with `DOUBLE`

precision. You should try to avoid using unsigned long long values if they resolve to be larger than 63 bits (9223372036854775807) for anything other than bit fields.

- You can have up to 255 `ENUM` and `SET` columns in one table.
- In `MIN()`, `MAX()`, and other aggregate functions, MySQL currently compares `ENUM` and `SET` columns by their string value rather than by the string's relative position in the set.
- In an `UPDATE` statement, columns are updated from left to right. If you refer to an updated column, you get the updated value instead of the original value. For example, the following statement increments `KEY` by `2`, **not** `1`:

```
mysql> UPDATE tbl_name SET KEY=KEY+1,KEY=KEY+1;
```

- You can refer to multiple temporary tables in the same query, but you cannot refer to any given temporary table more than once. For example, the following does not work:

```
mysql> SELECT * FROM temp_table, temp_table AS t2;
ERROR 1137: Can't reopen table: 'temp_table'
```

- The optimizer may handle `DISTINCT` differently when you are using “hidden” columns in a join than when you are not. In a join, hidden columns are counted as part of the result (even if they are not shown), whereas in normal queries, hidden columns do not participate in the `DISTINCT` comparison.

  An example of this is:

```
SELECT DISTINCT mp3id FROM band_downloads
       WHERE userid = 9 ORDER BY id DESC;
```

  and

```
SELECT DISTINCT band_downloads.mp3id
       FROM band_downloads,band_mp3
       WHERE band_downloads.userid = 9
       AND band_mp3.id = band_downloads.mp3id
       ORDER BY band_downloads.id DESC;
```

  In the second case, you may get two identical rows in the result set (because the values in the hidden `id` column may differ).

  Note that this happens only for queries that do not have the `ORDER BY` columns in the result.

- If you execute a `PROCEDURE` on a query that returns an empty set, in some cases the `PROCEDURE` does not transform the columns.
- Creation of a table of type `MERGE` does not check whether the underlying tables are compatible types.
- If you use `ALTER TABLE` to add a `UNIQUE` index to a table used in a `MERGE` table and then add a normal index on the `MERGE` table, the key order is different for the tables if there was an old, non-`UNIQUE` key in the table. This is because `ALTER TABLE` puts `UNIQUE` indexes before normal indexes to be able to detect duplicate keys as early as possible.

Appendix C Restrictions and Limits

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The discussion here describes restrictions that apply to the use of MySQL features such as subqueries or views.

C.1 Restrictions on Stored Programs

These restrictions apply to the features described in Chapter 23, *Stored Programs and Views*.

Some of the restrictions noted here apply to all stored routines; that is, both to stored procedures and stored functions. There are also some restrictions specific to stored functions but not to stored procedures.

The restrictions for stored functions also apply to triggers. There are also some restrictions specific to triggers.

The restrictions for stored procedures also apply to the `DO` clause of Event Scheduler event definitions. There are also some restrictions specific to events.

SQL Statements Not Permitted in Stored Routines

Stored routines cannot contain arbitrary SQL statements. The following statements are not permitted:

- The locking statements `LOCK TABLES` and `UNLOCK TABLES`.
- `ALTER VIEW`.
- `LOAD DATA` and `LOAD TABLE`.
- SQL prepared statements (`PREPARE`, `EXECUTE`, `DEALLOCATE PREPARE`) can be used in stored procedures, but not stored functions or triggers. Thus, stored functions and triggers cannot use dynamic SQL (where you construct statements as strings and then execute them).
- Generally, statements not permitted in SQL prepared statements are also not permitted in stored programs. For a list of statements supported as prepared statements, see Section 13.5, “Prepared SQL Statement Syntax”. Exceptions are `SIGNAL`, `RESIGNAL`, and `GET DIAGNOSTICS`, which are not permissible as prepared statements but are permitted in stored programs.
- Because local variables are in scope only during stored program execution, references to them are not permitted in prepared statements created within a stored program. Prepared statement scope is the current session, not the stored program, so the statement could be executed after the

program ends, at which point the variables would no longer be in scope. For example, `SELECT ... INTO local_var` cannot be used as a prepared statement. This restriction also applies to stored procedure and function parameters. See Section 13.5.1, “PREPARE Syntax”.

- Within all stored programs (stored procedures and functions, triggers, and events), the parser treats `BEGIN [WORK]` as the beginning of a `BEGIN ... END` block. To begin a transaction in this context, use `START TRANSACTION` instead.

Restrictions for Stored Functions

The following additional statements or operations are not permitted within stored functions. They are permitted within stored procedures, except stored procedures that are invoked from within a stored function or trigger. For example, if you use `FLUSH` in a stored procedure, that stored procedure cannot be called from a stored function or trigger.

- Statements that perform explicit or implicit commit or rollback. Support for these statements is not required by the SQL standard, which states that each DBMS vendor may decide whether to permit them.
- Statements that return a result set. This includes `SELECT` statements that do not have an `INTO var_list` clause and other statements such as `SHOW`, `EXPLAIN`, and `CHECK TABLE`. A function can process a result set either with `SELECT ... INTO var_list` or by using a cursor and `FETCH` statements. See Section 13.2.10.1, “SELECT ... INTO Syntax”, and Section 13.6.6, “Cursors”.
- `FLUSH` statements.
- Stored functions cannot be used recursively.
- A stored function or trigger cannot modify a table that is already being used (for reading or writing) by the statement that invoked the function or trigger.
- If you refer to a temporary table multiple times in a stored function under different aliases, a `Can't reopen table: 'tbl_name'` error occurs, even if the references occur in different statements within the function.
- `HANDLER ... READ` statements that invoke stored functions can cause replication errors and are disallowed.

Restrictions for Triggers

For triggers, the following additional restrictions apply:

- Triggers are not activated by foreign key actions.
- When using row-based replication, triggers on the slave are not activated by statements originating on the master. The triggers on the slave are activated when using statement-based replication. For more information, see Section 18.4.1.35, “Replication and Triggers”.
- The `RETURN` statement is not permitted in triggers, which cannot return a value. To exit a trigger immediately, use the `LEAVE` statement.
- Triggers are not permitted on tables in the `mysql` database. Nor are they permitted on `INFORMATION_SCHEMA` or `performance_schema` tables. Those tables are actually views and triggers are not permitted on views.
- The trigger cache does not detect when metadata of the underlying objects has changed. If a trigger uses a table and the table has changed since the trigger was loaded into the cache, the trigger operates using the outdated metadata.

Name Conflicts within Stored Routines

The same identifier might be used for a routine parameter, a local variable, and a table column. Also, the same local variable name can be used in nested blocks. For example:

```
CREATE PROCEDURE p (i INT)
BEGIN
  DECLARE i INT DEFAULT 0;
  SELECT i FROM t;
  BEGIN
    DECLARE i INT DEFAULT 1;
    SELECT i FROM t;
  END;
END;
```

In such cases, the identifier is ambiguous and the following precedence rules apply:

- A local variable takes precedence over a routine parameter or table column.
- A routine parameter takes precedence over a table column.
- A local variable in an inner block takes precedence over a local variable in an outer block.

The behavior that variables take precedence over table columns is nonstandard.

Replication Considerations

Use of stored routines can cause replication problems. This issue is discussed further in Section 23.7, “Binary Logging of Stored Programs”.

The `--replicate-wild-do-table=db_name.tbl_name` option applies to tables, views, and triggers. It does not apply to stored procedures and functions, or events. To filter statements operating on the latter objects, use one or more of the `--replicate-*-db` options.

Debugging Considerations

There are no stored routine debugging facilities.

Unsupported Syntax from the SQL:2003 Standard

The MySQL stored routine syntax is based on the SQL:2003 standard. The following items from that standard are not currently supported:

- `UNDO` handlers
- `FOR` loops

Concurrency Considerations

To prevent problems of interaction between sessions, when a client issues a statement, the server uses a snapshot of routines and triggers available for execution of the statement. That is, the server calculates a list of procedures, functions, and triggers that may be used during execution of the statement, loads them, and then proceeds to execute the statement. While the statement executes, it does not see changes to routines performed by other sessions.

For maximum concurrency, stored functions should minimize their side-effects; in particular, updating a table within a stored function can reduce concurrent operations on that table. A stored function acquires table locks before executing, to avoid inconsistency in the binary log due to mismatch of the order in which statements execute and when they appear in the log. When statement-based binary logging is used, statements that invoke a function are recorded rather than the statements executed within the function. Consequently, stored functions that update the same underlying tables do not execute in parallel. In contrast, stored procedures do not acquire table-level locks. All statements executed within stored procedures are written to the binary log, even for statement-based binary logging. See Section 23.7, “Binary Logging of Stored Programs”.

Event Scheduler Restrictions

The following limitations are specific to the Event Scheduler:

- Event names are handled in case-insensitive fashion. For example, you cannot have two events in the same database with the names `anEvent` and `AnEvent`.
- An event may not be created, altered, or dropped by a stored routine, trigger, or another event. An event also may not create, alter, or drop stored routines or triggers. (Bug #16409, Bug #18896)
- DDL statements on events are prohibited while a `LOCK TABLES` statement is in effect.
- Event timings using the intervals `YEAR`, `QUARTER`, `MONTH`, and `YEAR_MONTH` are resolved in months; those using any other interval are resolved in seconds. There is no way to cause events scheduled to occur at the same second to execute in a given order. In addition—due to rounding, the nature of threaded applications, and the fact that a nonzero length of time is required to create events and to signal their execution—events may be delayed by as much as 1 or 2 seconds. However, the time shown in the `INFORMATION_SCHEMA.EVENTS` table's `LAST_EXECUTED` column or the `mysql.event` table's `last_executed` column is always accurate to within one second of the actual event execution time. (See also Bug #16522.)
- Each execution of the statements contained in the body of an event takes place in a new connection; thus, these statements has no effect in a given user session on the server's statement counts such as `Com_select` and `Com_insert` that are displayed by `SHOW STATUS`. However, such counts *are* updated in the global scope. (Bug #16422)
- Events do not support times later than the end of the Unix Epoch; this is approximately the beginning of the year 2038. Such dates are specifically not permitted by the Event Scheduler. (Bug #16396)
- References to stored functions, user-defined functions, and tables in the `ON SCHEDULE` clauses of `CREATE EVENT` and `ALTER EVENT` statements are not supported. These sorts of references are not permitted. (See Bug #22830 for more information.)

C.2 Restrictions on Condition Handling

`SIGNAL`, `RESIGNAL`, and `GET DIAGNOSTICS` are not permissible as prepared statements. For example, this statement is invalid:

```
PREPARE stmt1 FROM 'SIGNAL SQLSTATE "02000"';
```

`SQLSTATE` values in class `'04'` are not treated specially. They are handled the same as other exceptions.

Standard SQL has a diagnostics area stack, containing a diagnostics area for each nested execution context. Standard SQL syntax includes `GET STACKED DIAGNOSTICS` for referring to stacked areas. MySQL does not support the `STACKED` keyword because there is a single diagnostics area containing information from the most recent statement that wrote to it. See also Section 13.6.7.7, “The MySQL Diagnostics Area”.

In standard SQL, the first condition relates to the `SQLSTATE` value returned for the previous SQL statement. In MySQL, this is not guaranteed, so to get the main error, you cannot do this:

```
GET DIAGNOSTICS CONDITION 1 @errno = MYSQL_ERRNO;
```

Instead, do this:

```
GET DIAGNOSTICS @cno = NUMBER;
GET DIAGNOSTICS CONDITION @cno @errno = MYSQL_ERRNO;
```

C.3 Restrictions on Server-Side Cursors

Server-side cursors are implemented in the C API using the `mysql_stmt_attr_set()` function. The same implementation is used for cursors in stored routines. A server-side cursor enables a result set to be generated on the server side, but not transferred to the client except for those rows that the client

requests. For example, if a client executes a query but is only interested in the first row, the remaining rows are not transferred.

In MySQL, a server-side cursor is materialized into an internal temporary table. Initially, this is a `MEMORY` table, but is converted to a `MyISAM` table when its size exceeds the minimum value of the `max_heap_table_size` and `tmp_table_size` system variables. The same restrictions apply to internal temporary tables created to hold the result set for a cursor as for other uses of internal temporary tables. See Section 8.4.4, “Internal Temporary Table Use in MySQL”. One limitation of the implementation is that for a large result set, retrieving its rows through a cursor might be slow.

Cursors are read only; you cannot use a cursor to update rows.

`UPDATE WHERE CURRENT OF` and `DELETE WHERE CURRENT OF` are not implemented, because updatable cursors are not supported.

Cursors are nonholdable (not held open after a commit).

Cursors are asensitive.

Cursors are nonscrollable.

Cursors are not named. The statement handler acts as the cursor ID.

You can have open only a single cursor per prepared statement. If you need several cursors, you must prepare several statements.

You cannot use a cursor for a statement that generates a result set if the statement is not supported in prepared mode. This includes statements such as `CHECK TABLE`, `HANDLER READ`, and `SHOW BINLOG EVENTS`.

C.4 Restrictions on Subqueries

- In general, you cannot modify a table and select from the same table in a subquery. For example, this limitation applies to statements of the following forms:

```
DELETE FROM t WHERE ... (SELECT ... FROM t ...);
UPDATE t ... WHERE col = (SELECT ... FROM t ...);
{INSERT|REPLACE} INTO t (SELECT ... FROM t ...);
```

  Exception: The preceding prohibition does not apply if for the modified table you are using a derived table (subquery in the `FROM` clause) and that derived table is materialized rather than merged into the outer query. (See Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.) Example:

```
UPDATE t ... WHERE col = (SELECT * FROM (SELECT ... FROM t...) AS dt ...);
```

  Here the result from the derived table is materialized as a temporary table, so the relevant rows in `t` have already been selected by the time the update to `t` takes place.

  In general, you may be able to influence the optimizer to materialize a derived table by adding a `NO_MERGE` optimizer hint. See Section 8.9.2, “Optimizer Hints”.

- Row comparison operations are only partially supported:

  - For *`expr [NOT] IN subquery`*, *`expr`* can be an *n*-tuple (specified using row constructor syntax) and the subquery can return rows of *n*-tuples. The permitted syntax is therefore more specifically expressed as *`row_constructor [NOT] IN table_subquery`*

  - For *`expr op {ALL|ANY|SOME} subquery`*, *`expr`* must be a scalar value and the subquery must be a column subquery; it cannot return multiple-column rows.

  In other words, for a subquery that returns rows of *n*-tuples, this is supported:

```
(expr_1, ..., expr_n) [NOT] IN table_subquery
```

But this is not supported:

```
(expr_1, ..., expr_n) op {ALL|ANY|SOME} subquery
```

The reason for supporting row comparisons for `IN` but not for the others is that `IN` is implemented by rewriting it as a sequence of `=` comparisons and `AND` operations. This approach cannot be used for `ALL`, `ANY`, or `SOME`.

- Subqueries in the `FROM` clause cannot be correlated subqueries. They are materialized in whole (evaluated to produce a result set) during query execution, so they cannot be evaluated per row of the outer query. The optimizer delays materialization until the result is needed, which may permit materialization to be avoided. See Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”.

- MySQL does not support `LIMIT` in subqueries for certain subquery operators:

```
mysql> SELECT * FROM t1
    ->   WHERE s1 IN (SELECT s2 FROM t2 ORDER BY s1 LIMIT 1);
ERROR 1235 (42000): This version of MySQL doesn't yet support
 'LIMIT & IN/ALL/ANY/SOME subquery'
```

- MySQL permits a subquery to refer to a stored function that has data-modifying side effects such as inserting rows into a table. For example, if `f()` inserts rows, the following query can modify data:

```
SELECT ... WHERE x IN (SELECT f() ...);
```

This behavior is an extension to the SQL standard. In MySQL, it can produce indeterminate results because `f()` might be executed a different number of times for different executions of a given query depending on how the optimizer chooses to handle it.

For statement-based or mixed-format replication, one implication of this indeterminism is that such a query can produce different results on the master and its slaves.

C.5 Restrictions on Views

View processing is not optimized:

- It is not possible to create an index on a view.

- Indexes can be used for views processed using the merge algorithm. However, a view that is processed with the temptable algorithm is unable to take advantage of indexes on its underlying tables (although indexes can be used during generation of the temporary tables).

There is a general principle that you cannot modify a table and select from the same table in a subquery. See Section C.4, “Restrictions on Subqueries”.

The same principle also applies if you select from a view that selects from the table, if the view selects from the table in a subquery and the view is evaluated using the merge algorithm. Example:

```
CREATE VIEW v1 AS
SELECT * FROM t2 WHERE EXISTS (SELECT 1 FROM t1 WHERE t1.a = t2.a);

UPDATE t1, v2 SET t1.a = 1 WHERE t1.b = v2.b;
```

If the view is evaluated using a temporary table, you *can* select from the table in the view subquery and still modify that table in the outer query. In this case the view will be stored in a temporary table and thus you are not really selecting from the table in a subquery and modifying it “at the same time.”

(This is another reason you might wish to force MySQL to use the temptable algorithm by specifying `ALGORITHM = TEMPTABLE` in the view definition.)

You can use `DROP TABLE` or `ALTER TABLE` to drop or alter a table that is used in a view definition. No warning results from the `DROP` or `ALTER` operation, even though this invalidates the view. Instead, an error occurs later, when the view is used. `CHECK TABLE` can be used to check for views that have been invalidated by `DROP` or `ALTER` operations.

With regard to view updatability, the overall goal for views is that if any view is theoretically updatable, it should be updatable in practice. MySQL as quickly as possible. Many theoretically updatable views can be updated now, but limitations still exist. For details, see Section 23.5.3, “Updatable and Insertable Views”.

There exists a shortcoming with the current implementation of views. If a user is granted the basic privileges necessary to create a view (the `CREATE VIEW` and `SELECT` privileges), that user will be unable to call `SHOW CREATE VIEW` on that object unless the user is also granted the `SHOW VIEW` privilege.

That shortcoming can lead to problems backing up a database with `mysqldump`, which may fail due to insufficient privileges. This problem is described in Bug #22062.

The workaround to the problem is for the administrator to manually grant the `SHOW VIEW` privilege to users who are granted `CREATE VIEW`, since MySQL doesn't grant it implicitly when views are created.

Views do not have indexes, so index hints do not apply. Use of index hints when selecting from a view is not permitted.

`SHOW CREATE VIEW` displays view definitions using an `AS` *`alias_name`* clause for each column. If a column is created from an expression, the default alias is the expression text, which can be quite long. Aliases for column names in `CREATE VIEW` statements are checked against the maximum column length of 64 characters (not the maximum alias length of 256 characters). As a result, views created from the output of `SHOW CREATE VIEW` fail if any column alias exceeds 64 characters. This can cause problems in the following circumstances for views with too-long aliases:

- View definitions fail to replicate to newer slaves that enforce the column-length restriction.
- Dump files created with `mysqldump` cannot be loaded into servers that enforce the column-length restriction.

A workaround for either problem is to modify each problematic view definition to use aliases that provide shorter column names. Then the view will replicate properly, and can be dumped and reloaded without causing an error. To modify the definition, drop and create the view again with `DROP VIEW` and `CREATE VIEW`, or replace the definition with `CREATE OR REPLACE VIEW`.

For problems that occur when reloading view definitions in dump files, another workaround is to edit the dump file to modify its `CREATE VIEW` statements. However, this does not change the original view definitions, which may cause problems for subsequent dump operations.

C.6 Restrictions on XA Transactions

XA transaction support is limited to the `InnoDB` storage engine.

For “external XA,” a MySQL server acts as a Resource Manager and client programs act as Transaction Managers. For “Internal XA”, storage engines within a MySQL server act as RMs, and the server itself acts as a TM. Internal XA support is limited by the capabilities of individual storage engines. Internal XA is required for handling XA transactions that involve more than one storage engine. The implementation of internal XA requires that a storage engine support two-phase commit at the table handler level, and currently this is true only for `InnoDB`.

For `XA START`, the `JOIN` and `RESUME` clauses are not supported.

For `XA END`, the `SUSPEND [FOR MIGRATE]` clause is not supported.

The requirement that the *`bqual`* part of the *`xid`* value be different for each XA transaction within a global transaction is a limitation of the current MySQL XA implementation. It is not part of the XA specification.

An XA transaction is written to the binary log in two parts. When `XA PREPARE` is issued, the first part of the transaction up to `XA PREPARE` is written using an initial GTID. A `XA_prepare_log_event` is used to identify such transactions in the binary log. When `XA COMMIT` or `XA ROLLBACK` is issued, a second part of the transaction containing only the `XA COMMIT` or `XA ROLLBACK` statement is written using a second GTID. Note that the initial part of the transaction, identified by `XA_prepare_log_event`, is not necessarily followed by its `XA COMMIT` or `XA ROLLBACK`, which can cause interleaved binary logging of any two XA transactions. The two parts of the XA transaction can even appear in different binary log files. This means that an XA transaction in `PREPARED` state is now persistent until an explicit `XA COMMIT` or `XA ROLLBACK` statement is issued, ensuring that XA transactions are compatible with replication.

The following restrictions exist for using XA transactions:

- XA is not fully crash-safe with respect to the binary log (on the master). If there is a crash before `XA PREPARE`, between `XA PREPARE` and `XA COMMIT` (or `XA ROLLBACK`), or after `XA COMMIT` (or `XA ROLLBACK`), the server and binary log are correctly recovered and taken to a consistent state. However, if there is a crash in the middle of the execution of one of these statements, the server may not be able to recover to a correct state, leaving the server state and the binary log in an inconsistent state.
- XA does not work with `relay-log-info-repository=TABLE`.
- XA does not work with replication filters or binary log filters. Filters are permitted as long as they do not render any XA transactions empty. Filters that filter out XA transactions may cause the slave to stop with an error.
- In case GTIDs are enabled and the slave does not use either `log-bin=OFF` or does not use `log-slave-updates`, XA transactions are not crash-safe with respect to GTIDs on the slave. If the slave stops unexpectedly while applying an `XA PREPARE` or `XA COMMIT`, then after recovery `@@GLOBAL.GTID_EXECUTED` may not correctly describe the transactions that have been applied on the slave.
- `FLUSH TABLES WITH READ LOCK` is not compatible with XA transactions.

Note

Prior to MySQL 5.7.7, XA transactions were not compatible with replication. This was because an XA transaction that was in `PREPARED` state would be rolled back on clean server shutdown or client disconnect. Similarly, an XA transaction that was in `PREPARED` state would still exist in `PREPARED` state in case the server was shutdown abnormally and then started again, but the contents of the transaction could not be written to the binary log. In both of these situations the XA transaction could not be replicated correctly.

C.7 Restrictions on Character Sets

- Identifiers are stored in `mysql` database tables (`user`, `db`, and so forth) using `utf8`, but identifiers can contain only characters in the Basic Multilingual Plane (BMP). Supplementary characters are not permitted in identifiers.
- The `ucs2`, `utf16`, `utf16le`, and `utf32` character sets have the following restrictions:
  - They cannot be used as a client character set, which means that they do not work for `SET NAMES` or `SET CHARACTER SET`. (See Section 10.1.4, “Connection Character Sets and Collations”.)
  - It is currently not possible to use `LOAD DATA INFILE` to load data files that use these character sets.

  - `FULLTEXT` indexes cannot be created on a column that uses any of these character sets. However, you can perform `IN BOOLEAN MODE` searches on the column without an index.
  - The use of `ENCRYPT()` with these character sets is not recommended because the underlying system call expects a string terminated by a zero byte.
- The `REGEXP` and `RLIKE` operators work in byte-wise fashion, so they are not multibyte safe and may produce unexpected results with multibyte character sets. In addition, these operators compare characters by their byte values and accented characters may not compare as equal even if a given collation treats them as equal.

C.8 Restrictions on Performance Schema

The Performance Schema avoids using mutexes to collect or produce data, so there are no guarantees of consistency and results can sometimes be incorrect. Event values in `performance_schema` tables are nondeterministic and nonrepeatable.

If you save event information in another table, you should not assume that the original events will still be available later. For example, if you select events from a `performance_schema` table into a temporary table, intending to join that table with the original table later, there might be no matches.

`mysqldump` and `BACKUP DATABASE` ignore tables in the `performance_schema` database.

Tables in the `performance_schema` database cannot be locked with `LOCK TABLES`, except the `setup_xxx` tables.

Tables in the `performance_schema` database cannot be indexed.

Results for queries that refer to tables in the `performance_schema` database are not saved in the query cache.

Tables in the `performance_schema` database are not replicated.

The types of timers might vary per platform. The `performance_timers` table shows which event timers are available. If the values in this table for a given timer name are `NULL`, that timer is not supported on your platform.

Instruments that apply to storage engines might not be implemented for all storage engines. Instrumentation of each third-party engine is the responsibility of the engine maintainer.

C.9 Restrictions on Pluggable Authentication

The first part of this section describes general restrictions on the applicability of the pluggable authentication framework described at Section 6.3.10, “Pluggable Authentication”. The second part describes how third-party connector developers can determine the extent to which a connector can take advantage of pluggable authentication capabilities and what steps to take to become more compliant.

The term “native authentication” used here refers to authentication against passwords stored in the `mysql.user` table. This is the same authentication method provided by older MySQL servers, before pluggable authentication was implemented. It remains the default method, although now it is implemented using plugins. “Windows native authentication” refers to authentication using the credentials of a user who has already logged in to Windows, as implemented by the Windows Native Authentication plugin (“Windows plugin” for short).

General Pluggable Authentication Restrictions

- **Connector/C, Connector/C++:** Clients that use these connectors can connect to the server only through accounts that use native authentication.

Exception: A connector supports pluggable authentication if it was built to link to `libmysqlclient` dynamically (rather than statically) and it loads the current version of `libmysqlclient` if that version is installed, or if the connector is recompiled from source to link against the current `libmysqlclient`.

- **Connector/Net:** Clients that use Connector/Net can connect to the server through accounts that use native authentication or the Windows plugin.
- **Connector/PHP:** Clients that use this connector can connect to the server only through accounts that use native authentication, when compiled using the MySQL native driver for PHP (`mysqlnd`).
- **Windows native authentication:** Connecting through an account that uses the Windows plugin requires Windows Domain setup. Without it, NTLM authentication is used and then only local connections are possible; that is, the client and server must run on the same computer.
- **Proxy users:** Proxy user support is available to the extent that clients can connect through accounts authenticated with plugins that implement proxy user capability (that is, plugins that can return a user name different from that of the connecting user). For example, the native authentication plugins do not support proxy users, whereas the PAM and Windows plugins do.
- **Replication**: Replication slaves can employ not only master accounts using native authentication, but can also connect through master accounts that use nonnative authentication if the required client-side plugin is available. If the plugin is built into `libmysqlclient`, it is available by default. Otherwise, the plugin must be installed on the slave side in the directory named by the slave `plugin_dir` system variable.
- **`FEDERATED` tables:** A `FEDERATED` table can access the remote table only through accounts on the remote server that use native authentication.

Pluggable Authentication and Third-Party Connectors

Third-party connector developers can use the following guidelines to determine readiness of a connector to take advantage of pluggable authentication capabilities and what steps to take to become more compliant:

- An existing connector to which no changes have been made uses native authentication and clients that use the connector can connect to the server only through accounts that use native authentication. *However, you should test the connector against a recent version of the server to verify that such connections still work without problem.*

  Exception: A connector might work with pluggable authentication without any changes if it links to `libmysqlclient` dynamically (rather than statically) and it loads the current version of `libmysqlclient` if that version is installed.

- To take advantage of pluggable authentication capabilities, a connector that is `libmysqlclient`-based should be relinked against the current version of `libmysqlclient`. This enables the connector to support connections though accounts that require client-side plugins now built into `libmysqlclient` (such as the cleartext plugin needed for PAM authentication and the Windows plugin needed for Windows native authentication). Linking with a current `libmysqlclient` also enables the connector to access client-side plugins installed in the default MySQL plugin directory (typically the directory named by the default value of the local server's `plugin_dir` system variable).

  If a connector links to `libmysqlclient` dynamically, it must be ensured that the newer version of `libmysqlclient` is installed on the client host and that the connector loads it at runtime.

- Another way for a connector to support a given authentication method is to implement it directly in the client/server protocol. Connector/Net uses this approach to provide support for Windows native authentication.

- If a connector should be able to load client-side plugins from a directory different from the default plugin directory, it must implement some means for client users to specify the directory. Possibilities for this include a command-line option or environment variable from which the connector can obtain the directory name. Standard MySQL client programs such as `mysql` and `mysqladmin` implement a `--plugin-dir` option. See also Section 27.7.13, “C API Client Plugin Functions”.
- Proxy user support by a connector depends, as described earlier in this section, on whether the authentication methods that it supports permit proxy users.

C.10 Limits in MySQL

This section lists current limits in MySQL 8.0.

C.10.1 Limits on Joins

The maximum number of tables that can be referenced in a single join is 61. This includes a join handled by merging derived tables (subqueries) and views in the `FROM` clause into the outer query block (see Section 8.2.2.3, “Optimizing Derived Tables, View References, and Common Table Expressions”). It also applies to the number of tables that can be referenced in the definition of a view.

C.10.2 Limits on Number of Databases and Tables

MySQL has no limit on the number of databases. The underlying file system may have a limit on the number of directories.

MySQL has no limit on the number of tables. The underlying file system may have a limit on the number of files that represent tables. Individual storage engines may impose engine-specific constraints. `InnoDB` permits up to 4 billion tables.

C.10.3 Limits on Table Size

The effective maximum table size for MySQL databases is usually determined by operating system constraints on file sizes, not by MySQL internal limits. For up-to-date information operating system file size limits, refer to the documentation specific to your operating system.

Windows users, please note that FAT and VFAT (FAT32) are *not* considered suitable for production use with MySQL. Use NTFS instead.

If you encounter a full-table error, there are several reasons why it might have occurred:

- The disk might be full.
- You are using `InnoDB` tables and have run out of room in an `InnoDB` tablespace file. The maximum tablespace size is also the maximum size for a table. For tablespace size limits, see Section 15.8.1.7, “Limits on InnoDB Tables”.

  Generally, partitioning of tables into multiple tablespace files is recommended for tables larger than 1TB in size.
- You have hit an operating system file size limit. For example, you are using `MyISAM` tables on an operating system that supports files only up to 2GB in size and you have hit this limit for the data file or index file.
- You are using a `MyISAM` table and the space required for the table exceeds what is permitted by the internal pointer size. `MyISAM` permits data and index files to grow up to 256TB by default, but this limit can be changed up to the maximum permissible size of 65,536TB ($256^7 - 1$ bytes).

  If you need a `MyISAM` table that is larger than the default limit and your operating system supports large files, the `CREATE TABLE` statement supports `AVG_ROW_LENGTH` and `MAX_ROWS` options. See Section 13.1.15, “CREATE TABLE Syntax”. The server uses these options to determine how large a table to permit.

If the pointer size is too small for an existing table, you can change the options with `ALTER TABLE` to increase a table's maximum permissible size. See Section 13.1.7, “ALTER TABLE Syntax”.

```
ALTER TABLE tbl_name MAX_ROWS=1000000000 AVG_ROW_LENGTH=nnn;
```

You have to specify `AVG_ROW_LENGTH` only for tables with `BLOB` or `TEXT` columns; in this case, MySQL can't optimize the space required based only on the number of rows.

To change the default size limit for `MyISAM` tables, set the `myisam_data_pointer_size`, which sets the number of bytes used for internal row pointers. The value is used to set the pointer size for new tables if you do not specify the `MAX_ROWS` option. The value of `myisam_data_pointer_size` can be from 2 to 7. A value of 4 permits tables up to 4GB; a value of 6 permits tables up to 256TB.

You can check the maximum data and index sizes by using this statement:

```
SHOW TABLE STATUS FROM db_name LIKE 'tbl_name';
```

You also can use `myisamchk -dv /path/to/table-index-file`. See Section 13.7.5, “SHOW Syntax”, or Section 4.6.4, “`myisamchk` — MyISAM Table-Maintenance Utility”.

Other ways to work around file-size limits for `MyISAM` tables are as follows:

- If your large table is read only, you can use `myisampack` to compress it. `myisampack` usually compresses a table by at least 50%, so you can have, in effect, much bigger tables. `myisampack` also can merge multiple tables into a single table. See Section 4.6.6, “`myisampack` — Generate Compressed, Read-Only MyISAM Tables”.
- MySQL includes a `MERGE` library that enables you to handle a collection of `MyISAM` tables that have identical structure as a single `MERGE` table. See Section 16.7, “The MERGE Storage Engine”.

- You are using the `MEMORY` (`HEAP`) storage engine; in this case you need to increase the value of the `max_heap_table_size` system variable. See Section 5.1.5, “Server System Variables”.

C.10.4 Limits on Table Column Count and Row Size

Column Count Limits

MySQL has hard limit of 4096 columns per table, but the effective maximum may be less for a given table. The exact column limit depends on several factors:

- The maximum row size for a table constrains the number (and possibly size) of columns because the total length of all columns cannot exceed this size. See Row Size Limits.
- The storage requirements of individual columns constrain the number of columns that fit within a given maximum row size. Storage requirements for some data types depend on factors such as storage engine, storage format, and character set. See Section 11.8, “Data Type Storage Requirements”.
- Storage engines may impose additional restrictions that limit table column count. For example, `InnoDB` has a limit of 1017 columns per table. See Section 15.8.1.7, “Limits on InnoDB Tables”. For information about other storage engines, see Chapter 16, *Alternative Storage Engines*.

Row Size Limits

The maximum row size for a given table is determined by several factors:

- The internal representation of a MySQL table has a maximum row size limit of 65,535 bytes, even if the storage engine is capable of supporting larger rows. `BLOB` and `TEXT` columns only contribute 9 to 12 bytes toward the row size limit because their contents are stored separately from the rest of the row.

- The maximum row size for an InnoDB table, which applies to data stored locally within a database page, is slightly less than half a page for 4KB, 8KB, 16KB, and 32KB innodb_page_size settings. For example, the maximum row size is slightly less than 8KB for the default 16KB InnoDB page size. For 64KB pages, the maximum row size is slightly less than 16KB. See Section 15.8.1.7, “Limits on InnoDB Tables”.

  If a row containing variable-length columns exceeds the InnoDB maximum row size, InnoDB selects variable-length columns for external off-page storage until the row fits within the InnoDB row size limit. The amount of data stored locally for variable-length columns that are stored off-page differs by row format. For more information, see Section 15.10, “InnoDB Row Storage and Row Formats”.

- Different storage formats use different amounts of page header and trailer data, which affects the amount of storage available for rows.

  - For information about InnoDB row formats, see Section 15.10, “InnoDB Row Storage and Row Formats”, and Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

  - For information about MyISAM storage formats, see Section 16.2.3, “MyISAM Table Storage Formats”.

Row Size Limit Examples

- The MySQL maximum row size limit of 65,535 bytes is demonstrated in the following InnoDB and MyISAM examples. The limit is enforced regardless of storage engine, even though the storage engine may be capable of supporting larger rows.

```
mysql> CREATE TABLE t (a VARCHAR(10000), b VARCHAR(10000),
       c VARCHAR(10000), d VARCHAR(10000), e VARCHAR(10000),
       f VARCHAR(10000), g VARCHAR(6000)) ENGINE=InnoDB CHARACTER SET latin1;
ERROR 1118 (42000): Row size too large. The maximum row size for the used
table type, not counting BLOBs, is 65535. This includes storage overhead,
check the manual. You have to change some columns to TEXT or BLOBs
```

```
mysql> CREATE TABLE t (a VARCHAR(10000), b VARCHAR(10000),
       c VARCHAR(10000), d VARCHAR(10000), e VARCHAR(10000),
       f VARCHAR(10000), g VARCHAR(6000)) ENGINE=MyISAM CHARACTER SET latin1;
ERROR 1118 (42000): Row size too large. The maximum row size for the used
table type, not counting BLOBs, is 65535. This includes storage overhead,
check the manual. You have to change some columns to TEXT or BLOBs
```

  In the following MyISAM example, changing a column to TEXT avoids the 65,535-byte row size limit and permits the operation to succeed because BLOB and TEXT columns only contribute 9 to 12 bytes toward the row size.

```
mysql> CREATE TABLE t (a VARCHAR(10000), b VARCHAR(10000),
       c VARCHAR(10000), d VARCHAR(10000), e VARCHAR(10000),
       f VARCHAR(10000), g TEXT(6000)) ENGINE=MyISAM CHARACTER SET latin1;
Query OK, 0 rows affected (0.02 sec)
```

  The operation succeeds for an InnoDB table because changing a column to TEXT avoids the MySQL 65,535-byte row size limit, and InnoDB off-page storage of variable-length columns avoids the InnoDB row size limit.

```
mysql> CREATE TABLE t (a VARCHAR(10000), b VARCHAR(10000),
       c VARCHAR(10000), d VARCHAR(10000), e VARCHAR(10000),
       f VARCHAR(10000), g TEXT(6000)) ENGINE=InnoDB CHARACTER SET latin1;
Query OK, 0 rows affected (0.02 sec)
```

- Storage for variable-length columns includes length bytes, which are counted toward the row size. For example, a VARCHAR(255) CHARACTER SET utf8mb3 column takes two bytes to store the length of the value, so each value can take up to 767 bytes.

The statement to create table `t1` succeeds because the columns require 32,765 + 2 bytes and 32,766 + 2 bytes, which falls within the maximum row size of 65,535 bytes:

```
mysql> CREATE TABLE t1
       (c1 VARCHAR(32765) NOT NULL, c2 VARCHAR(32766) NOT NULL)
       ENGINE = InnoDB CHARACTER SET latin1;
Query OK, 0 rows affected (0.02 sec)
```

The statement to create table `t2` fails because, although the column length is within the maximum length of 65,535 bytes, two additional bytes are required to record the length, which causes the row size to exceed 65,535 bytes:

```
mysql> CREATE TABLE t2
       (c1 VARCHAR(65535) NOT NULL)
       ENGINE = InnoDB CHARACTER SET latin1;
ERROR 1118 (42000): Row size too large. The maximum row size for the used
table type, not counting BLOBs, is 65535. This includes storage overhead,
check the manual. You have to change some columns to TEXT or BLOBs
```

Reducing the column length to 65,533 or less permits the statement to succeed.

```
mysql> CREATE TABLE t2
       (c1 VARCHAR(65533) NOT NULL)
       ENGINE = InnoDB CHARACTER SET latin1;
Query OK, 0 rows affected (0.01 sec)
```

- For `MyISAM` tables, `NULL` columns require additional space in the row to record whether their values are `NULL`. Each `NULL` column takes one bit extra, rounded up to the nearest byte.

  The statement to create table `t3` fails because `MyISAM` requires space for `NULL` columns in addition to the space required for variable-length column length bytes, causing the row size to exceed 65,535 bytes:

  ```
  mysql> CREATE TABLE t3
         (c1 VARCHAR(32765) NULL, c2 VARCHAR(32766) NULL)
         ENGINE = MyISAM CHARACTER SET latin1;
  ERROR 1118 (42000): Row size too large. The maximum row size for the used
  table type, not counting BLOBs, is 65535. This includes storage overhead,
  check the manual. You have to change some columns to TEXT or BLOBs
  ```

  For information about `InnoDB` `NULL` column storage, see Section 15.8.1.2, “The Physical Row Structure of an InnoDB Table”.

- `InnoDB` restricts row size (for data stored locally within the database page) to slightly less than half a database page for 4KB, 8KB, 16KB, and 32KB `innodb_page_size` settings, and to slightly less than 16KB for 64KB pages.

  The statement to create table `t4` fails because the defined columns exceed the row size limit for a 16KB `InnoDB` page.

  ```
  mysql> CREATE TABLE t4 (
         c1 CHAR(255),c2 CHAR(255),c3 CHAR(255),
         c4 CHAR(255),c5 CHAR(255),c6 CHAR(255),
         c7 CHAR(255),c8 CHAR(255),c9 CHAR(255),
         c10 CHAR(255),c11 CHAR(255),c12 CHAR(255),
         c13 CHAR(255),c14 CHAR(255),c15 CHAR(255),
         c16 CHAR(255),c17 CHAR(255),c18 CHAR(255),
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         c22 CHAR(255),c23 CHAR(255),c24 CHAR(255),
         c25 CHAR(255),c26 CHAR(255),c27 CHAR(255),
         c28 CHAR(255),c29 CHAR(255),c30 CHAR(255),
         c31 CHAR(255),c32 CHAR(255),c33 CHAR(255)
  ```

```
       ) ENGINE=InnoDB ROW_FORMAT=DYNAMIC DEFAULT CHARSET latin1;
ERROR 1118 (42000): Row size too large (> 8126). Changing some columns to TEXT or BLOB may help.
In current row format, BLOB prefix of 0 bytes is stored inline.
```

C.10.5 Windows Platform Limitations

The following limitations apply to use of MySQL on the Windows platform:

- **Process memory**

  On Windows 32-bit platforms, it is not possible by default to use more than 2GB of RAM within a single process, including MySQL. This is because the physical address limit on Windows 32-bit is 4GB and the default setting within Windows is to split the virtual address space between kernel (2GB) and user/applications (2GB).

  Some versions of Windows have a boot time setting to enable larger applications by reducing the kernel application. Alternatively, to use more than 2GB, use a 64-bit version of Windows.

- **File system aliases**

  When using `MyISAM` tables, you cannot use aliases within Windows link to the data files on another volume and then link back to the main MySQL `datadir` location.

  This facility is often used to move the data and index files to a RAID or other fast solution.

- **Limited number of ports**

  Windows systems have about 4,000 ports available for client connections, and after a connection on a port closes, it takes two to four minutes before the port can be reused. In situations where clients connect to and disconnect from the server at a high rate, it is possible for all available ports to be used up before closed ports become available again. If this happens, the MySQL server appears to be unresponsive even though it is running. Ports may be used by other applications running on the machine as well, in which case the number of ports available to MySQL is lower.

  For more information about this problem, see http://support.microsoft.com/default.aspx?scid=kb;en-us;196271.

- **`DATA DIRECTORY` and `INDEX DIRECTORY`**

  The `DATA DIRECTORY` option for `CREATE TABLE` is supported on Windows only for `InnoDB` tables, as described in Section 15.7.5, “Creating File-Per-Table Tablespaces Outside the Data Directory”. For `MyISAM` and other storage engines, the `DATA DIRECTORY` and `INDEX DIRECTORY` options for `CREATE TABLE` are ignored on Windows and any other platforms with a nonfunctional `realpath()` call.

- **`DROP DATABASE`**

  You cannot drop a database that is in use by another session.

- **Case-insensitive names**

  File names are not case sensitive on Windows, so MySQL database and table names are also not case sensitive on Windows. The only restriction is that database and table names must be specified using the same case throughout a given statement. See Section 9.2.2, “Identifier Case Sensitivity”.

- **Directory and file names**

  On Windows, MySQL Server supports only directory and file names that are compatible with the current ANSI code pages. For example, the following Japanese directory name will not work in the Western locale (code page 1252):

```
```

```
datadir="C:/私たちのプロジェクトのデータ"
```

The same limitation applies to directory and file names referred to in SQL statements, such as the data file path name in `LOAD DATA INFILE`.

- **The `\` path name separator character**

  Path name components in Windows are separated by the `\` character, which is also the escape character in MySQL. If you are using `LOAD DATA INFILE` or `SELECT ... INTO OUTFILE`, use Unix-style file names with `/` characters:

  ```
  mysql> LOAD DATA INFILE 'C:/tmp/skr.txt' INTO TABLE skr;
  mysql> SELECT * INTO OUTFILE 'C:/tmp/skr.txt' FROM skr;
  ```

  Alternatively, you must double the `\` character:

  ```
  mysql> LOAD DATA INFILE 'C:\\tmp\\skr.txt' INTO TABLE skr;
  mysql> SELECT * INTO OUTFILE 'C:\\tmp\\skr.txt' FROM skr;
  ```

- **Problems with pipes**

  Pipes do not work reliably from the Windows command-line prompt. If the pipe includes the character `^Z` / `CHAR(24)`, Windows thinks that it has encountered end-of-file and aborts the program.

  This is mainly a problem when you try to apply a binary log as follows:

  ```
  C:\> mysqlbinlog binary_log_file | mysql --user=root
  ```

  If you have a problem applying the log and suspect that it is because of a `^Z` / `CHAR(24)` character, you can use the following workaround:

  ```
  C:\> mysqlbinlog binary_log_file --result-file=/tmp/bin.sql
  C:\> mysql --user=root --execute "source /tmp/bin.sql"
  ```

  The latter command also can be used to reliably read in any SQL file that may contain binary data.

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SELECT * FROM t PARTITION ()

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SELECT * INTO OUTFILE 'file_name' FROM tbl_name

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SET NAMES 'charset_name'

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SET NAMES 'cp1251'

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SET NAMES charset_name

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SET NAMES default_character_set

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SET PASSWORD ... = PASSWORD()

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SET SESSION TRANSACTION

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SET sql_mode='modes'

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SET TIMESTAMP = value

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SET TRANSACTION

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SET var_name = value

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MySQL Glossary

These terms are commonly used in information about the MySQL database server. This glossary originated as a reference for terminology about the InnoDB storage engine, and the majority of definitions are InnoDB-related.

A

.ARM file
: Metadata for `ARCHIVE` tables. Contrast with **.ARZ file**. Files with this extension are always included in backups produced by the `mysqlbackup` command of the **MySQL Enterprise Backup** product.
See Also .ARZ file, MySQL Enterprise Backup, mysqlbackup command.

.ARZ file
: Data for ARCHIVE tables. Contrast with **.ARM file**. Files with this extension are always included in backups produced by the `mysqlbackup` command of the **MySQL Enterprise Backup** product.
See Also .ARM file, MySQL Enterprise Backup, mysqlbackup command.

ACID
: An acronym standing for atomicity, consistency, isolation, and durability. These properties are all desirable in a database system, and are all closely tied to the notion of a **transaction**. The transactional features of `InnoDB` adhere to the ACID principles.

  Transactions are **atomic** units of work that can be **committed** or **rolled back**. When a transaction makes multiple changes to the database, either all the changes succeed when the transaction is committed, or all the changes are undone when the transaction is rolled back.

  The database remains in a consistent state at all times — after each commit or rollback, and while transactions are in progress. If related data is being updated across multiple tables, queries see either all old values or all new values, not a mix of old and new values.

  Transactions are protected (isolated) from each other while they are in progress; they cannot interfere with each other or see each other's uncommitted data. This isolation is achieved through the **locking** mechanism. Experienced users can adjust the **isolation level**, trading off less protection in favor of increased performance and **concurrency**, when they can be sure that the transactions really do not interfere with each other.

  The results of transactions are durable: once a commit operation succeeds, the changes made by that transaction are safe from power failures, system crashes, race conditions, or other potential dangers that many non-database applications are vulnerable to. Durability typically involves writing to disk storage, with a certain amount of redundancy to protect against power failures or software crashes during write operations. (In `InnoDB`, the **doublewrite buffer** assists with durability.)
  See Also atomic, commit, concurrency, doublewrite buffer, isolation level, locking, rollback, transaction.

adaptive flushing
: An algorithm for **InnoDB** tables that smooths out the I/O overhead introduced by **checkpoints**. Instead of **flushing** all modified **pages** from the **buffer pool** to the **data files** at once, MySQL periodically flushes small sets of modified pages. The adaptive flushing algorithm extends this process by estimating the optimal rate to perform these periodic flushes, based on the rate of flushing and how fast **redo** information is generated.
See Also buffer pool, checkpoint, data files, flush, InnoDB, page, redo log.

adaptive hash index
: An optimization for `InnoDB` tables that can speed up lookups using `=` and `IN` operators, by constructing a **hash index** in memory. MySQL monitors index searches for `InnoDB` tables, and if queries could benefit from a hash index, it builds one automatically for index **pages** that are frequently accessed.
In a sense, the adaptive hash index configures MySQL at runtime to take advantage of ample main memory, coming closer to the architecture of main-memory databases. This feature is controlled by the `innodb_adaptive_hash_index` configuration option. Because this feature benefits some workloads and not others, and the memory used for the hash index is reserved in the **buffer pool**, typically you should benchmark with this feature both enabled and disabled.

The hash index is always built based on an existing `InnoDB` **secondary index**, which is organized as a **B-tree** structure. MySQL can build a hash index on a prefix of any length of the key defined for the B-tree, depending on the pattern of searches against the index. A hash index can be partial; the whole B-tree index does not need to be cached in the buffer pool.

In MySQL 5.6 and higher, another way to take advantage of fast single-value lookups with `InnoDB` tables is to use the `InnoDB` **memcached** plugin. See Section 15.19, “InnoDB memcached Plugin” for details.
See Also B-tree, buffer pool, hash index, memcached, page, secondary index.

AHI

Acronym for **adaptive hash index**.
See Also adaptive hash index.

AIO

Acronym for **asynchronous I/O**. You might see this acronym in `InnoDB` messages or keywords.
See Also asynchronous I/O.

application programming interface (API)

A set of functions or procedures. An API provides a stable set of names and types for functions, procedures, parameters, and return values.

apply

When a backup produced by the **MySQL Enterprise Backup** product does not include the most recent changes that occurred while the backup was underway, the process of updating the backup files to include those changes is known as the **apply** step. It is specified by the `apply-log` option of the `mysqlbackup` command.

Before the changes are applied, we refer to the files as a **raw backup**. After the changes are applied, we refer to the files as a **prepared backup**. The changes are recorded in the **ibbackup_logfile** file; once the apply step is finished, this file is no longer necessary.
See Also hot backup, ibbackup_logfile, MySQL Enterprise Backup, prepared backup, raw backup.

asynchronous I/O

A type of I/O operation that allows other processing to proceed before the I/O is completed. Also known as **non-blocking I/O** and abbreviated as **AIO**. `InnoDB` uses this type of I/O for certain operations that can run in parallel without affecting the reliability of the database, such as reading pages into the **buffer pool** that have not actually been requested, but might be needed soon.

Historically, `InnoDB` used asynchronous I/O on Windows systems only. Starting with the InnoDB Plugin 1.1 and MySQL 5.5, `InnoDB` uses asynchronous I/O on Linux systems. This change introduces a dependency on `libaio`. Asynchronous I/O on Linux systems is configured using the `innodb_use_native_aio` option, which is enabled by default. On other Unix-like systems, InnoDB uses synchronous I/O only.
See Also buffer pool, non-blocking I/O.

atomic

In the SQL context, **transactions** are units of work that either succeed entirely (when **committed**) or have no effect at all (when **rolled back**). The indivisible ("atomic") property of transactions is the “A” in the acronym **ACID**.
See Also ACID, commit, rollback, transaction.

atomic instruction

Special instructions provided by the CPU, to ensure that critical low-level operations cannot be interrupted.

auto-increment

A property of a table column (specified by the `AUTO_INCREMENT` keyword) that automatically adds an ascending sequence of values in the column. `InnoDB` supports auto-increment only for **primary key** columns.

It saves work for the developer, not to have to produce new unique values when inserting new rows. It provides useful information for the query optimizer, because the column is known to be not null and with unique values. The values from such a column can be used as lookup keys in various contexts, and because they are auto-generated there is no reason to ever change them; for this reason, primary key columns are often specified as auto-incrementing.

Auto-increment columns can be problematic with statement-based replication, because replaying the statements on a slave might not produce the same set of column values as on the master, due to timing issues. When you have an auto-incrementing primary key, you can use statement-based replication only with the setting `innodb_autoinc_lock_mode=1`. If you have `innodb_autoinc_lock_mode=2`, which allows higher concurrency for insert operations, use **row-based replication** rather than **statement-based replication**. The setting `innodb_autoinc_lock_mode=0` is the previous (traditional) default setting and should not be used except for compatibility purposes.
See Also auto-increment locking, innodb_autoinc_lock_mode, primary key, row-based replication, statement-based replication.

auto-increment locking
The convenience of an **auto-increment** primary key involves some tradeoff with concurrency. In the simplest case, if one transaction is inserting values into the table, any other transactions must wait to do their own inserts into that table, so that rows inserted by the first transaction receive consecutive primary key values. `InnoDB` includes optimizations, and the `innodb_autoinc_lock_mode` option, so that you can choose how to trade off between predictable sequences of auto-increment values and maximum **concurrency** for insert operations.
See Also auto-increment, concurrency, innodb_autoinc_lock_mode.

autocommit
A setting that causes a **commit** operation after each **SQL** statement. This mode is not recommended for working with `InnoDB` tables with **transactions** that span several statements. It can help performance for **read-only transactions** on `InnoDB` tables, where it minimizes overhead from **locking** and generation of **undo** data, especially in MySQL 5.6.4 and up. It is also appropriate for working with `MyISAM` tables, where transactions are not applicable.
See Also commit, locking, read-only transaction, SQL, transaction, undo.

availability
The ability to cope with, and if necessary recover from, failures on the host, including failures of MySQL, the operating system, or the hardware and maintenance activity that may otherwise cause downtime. Often paired with **scalability** as critical aspects of a large-scale deployment.
See Also scalability.

B

B-tree
A tree data structure that is popular for use in database indexes. The structure is kept sorted at all times, enabling fast lookup for exact matches (equals operator) and ranges (for example, greater than, less than, and `BETWEEN` operators). This type of index is available for most storage engines, such as `InnoDB` and `MyISAM`.

Because B-tree nodes can have many children, a B-tree is not the same as a binary tree, which is limited to 2 children per node.

Contrast with **hash index**, which is only available in the `MEMORY` storage engine. The `MEMORY` storage engine can also use B-tree indexes, and you should choose B-tree indexes for `MEMORY` tables if some queries use range operators.

The use of the term B-tree is intended as a reference to the general class of index design. B-tree structures used by MySQL storage engines may be regarded as variants due to sophistications not present in a classic B-tree design. For related information, refer to the `InnoDB` Page Structure Fil Header section of the MySQL Internals Manual.
See Also hash index.

backticks
Identifiers within MySQL SQL statements must be quoted using the backtick character (`) if they contain special characters or reserved words. For example, to refer to a table named `FOO#BAR` or a column named `SELECT`, you would specify the identifiers as `` `FOO#BAR` `` and `` `SELECT` ``. Since the backticks provide an extra level of safety, they are used extensively in program-generated SQL statements, where the identifier names might not be known in advance.

Many other database systems use double quotation marks (") around such special names. For portability, you can enable `ANSI_QUOTES` mode in MySQL and use double quotation marks instead of backticks to qualify identifier names.
See Also SQL.

backup
The process of copying some or all table data and metadata from a MySQL instance, for safekeeping. Can also refer to the set of copied files. This is a crucial task for DBAs. The reverse of this process is the **restore** operation.

With MySQL, **physical backups** are performed by the **MySQL Enterprise Backup** product, and **logical backups** are performed by the `mysqldump` command. These techniques have different characteristics in terms of size and representation of the backup data, and speed (especially speed of the restore operation).

Backups are further classified as **hot**, **warm**, or **cold** depending on how much they interfere with normal database operation. (Hot backups have the least interference, cold backups the most.)
See Also cold backup, hot backup, logical backup, MySQL Enterprise Backup, mysqldump, physical backup, warm backup.

base column
A non-generated table column upon which a stored generated column or virtual generated column is based. In other words, a base column is a non-generated table column that is part of a generated column definition.
See Also generated column, stored generated column, virtual generated column.

beta
An early stage in the life of a software product, when it is available only for evaluation, typically without a definite release number or a number less than 1. `InnoDB` does not use the beta designation, preferring an **early adopter** phase that can extend over several point releases, leading to a **GA** release.
See Also early adopter, GA.

binary log
A file containing a record of all statements that attempt to change table data. These statements can be replayed to bring slave servers up to date in a **replication** scenario, or to bring a database up to date after restoring table data from a backup. The binary logging feature can be turned on and off, although Oracle recommends always enabling it if you use replication or perform backups.

You can examine the contents of the binary log, or replay those statements during replication or recovery, by using the `mysqlbinlog` command. For full information about the binary log, see Section 5.4.4, “The Binary Log”. For MySQL configuration options related to the binary log, see Section 18.1.6.4, “Binary Logging Options and Variables”.

For the **MySQL Enterprise Backup** product, the file name of the binary log and the current position within the file are important details. To record this information for the master server when taking a backup in a replication context, you can specify the `--slave-info` option.

Prior to MySQL 5.0, a similar capability was available, known as the update log. In MySQL 5.0 and higher, the binary log replaces the update log.
See Also binlog, MySQL Enterprise Backup, replication.

binlog
An informal name for the **binary log** file. For example, you might see this abbreviation used in e-mail messages or forum discussions.
See Also binary log.

blind query expansion
A special mode of **full-text search** enabled by the `WITH QUERY EXPANSION` clause. It performs the search twice, where the search phrase for the second search is the original search phrase concatenated with the few most highly relevant documents from the first search. This technique is mainly applicable for short search phrases, perhaps only a single word. It can uncover relevant matches where the precise search term does not occur in the document.

See Also full-text search.

bottleneck
A portion of a system that is constrained in size or capacity, that has the effect of limiting overall throughput. For example, a memory area might be smaller than necessary; access to a single required resource might prevent multiple CPU cores from running simultaneously; or waiting for disk I/O to complete might prevent the CPU from running at full capacity. Removing bottlenecks tends to improve **concurrency**. For example, the ability to have multiple `InnoDB` **buffer pool** instances reduces contention when multiple sessions read from and write to the buffer pool simultaneously.
See Also buffer pool, concurrency.

bounce
A **shutdown** operation immediately followed by a restart. Ideally with a relatively short **warmup** period so that performance and throughput quickly return to a high level.
See Also shutdown.

buddy allocator
A mechanism for managing different-sized **pages** in the InnoDB **buffer pool**.
See Also buffer pool, page, page size.

buffer
A memory or disk area used for temporary storage. Data is buffered in memory so that it can be written to disk efficiently, with a few large I/O operations rather than many small ones. Data is buffered on disk for greater reliability, so that it can be recovered even when a **crash** or other failure occurs at the worst possible time. The main types of buffers used by InnoDB are the **buffer pool**, the **doublewrite buffer**, and the **change buffer**.
See Also buffer pool, change buffer, crash, doublewrite buffer.

buffer pool
The memory area that holds cached `InnoDB` data for both tables and indexes. For efficiency of high-volume read operations, the buffer pool is divided into **pages** that can potentially hold multiple rows. For efficiency of cache management, the buffer pool is implemented as a linked list of pages; data that is rarely used is aged out of the cache, using a variation of the **LRU** algorithm. On systems with large memory, you can improve concurrency by dividing the buffer pool into multiple **buffer pool instances**.

Several `InnoDB` status variables, `INFORMATION_SCHEMA` tables, and `performance_schema` tables help to monitor the internal workings of the buffer pool. Starting in MySQL 5.6, you can avoid a lengthy warmup period after restarting the server, particularly for instances with large buffer pools, by saving the buffer pool state at server shutdown and restoring the buffer pool to the same state at server startup. See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.
See Also buffer pool instance, LRU, page, warm up.

buffer pool instance
Any of the multiple regions into which the **buffer pool** can be divided, controlled by the `innodb_buffer_pool_instances` configuration option. The total memory size specified by `innodb_buffer_pool_size` is divided among all buffer pool instances. Typically, having multiple buffer pool instances is appropriate for systems that allocate multiple gigabytes to the `InnoDB` buffer pool, with each instance being one gigabyte or larger. On systems loading or looking up large amounts of data in the buffer pool from many concurrent sessions, having multiple buffer pool instances reduces contention for exclusive access to data structures that manage the buffer pool.
See Also buffer pool.

built-in
The built-in `InnoDB` storage engine within MySQL is the original form of distribution for the storage engine. Contrast with the **InnoDB Plugin**. Starting with MySQL 5.5, the InnoDB Plugin is merged back into the MySQL code base as the built-in `InnoDB` storage engine (known as InnoDB 1.1).

This distinction is important mainly in MySQL 5.1, where a feature or bug fix might apply to the InnoDB Plugin but not the built-in `InnoDB`, or vice versa.
See Also InnoDB.

business rules
: The relationships and sequences of actions that form the basis of business software, used to run a commercial company. Sometimes these rules are dictated by law, other times by company policy. Careful planning ensures that the relationships encoded and enforced by the database, and the actions performed through application logic, accurately reflect the real policies of the company and can handle real-life situations.

  For example, an employee leaving a company might trigger a sequence of actions from the human resources department. The human resources database might also need the flexibility to represent data about a person who has been hired, but not yet started work. Closing an account at an online service might result in data being removed from a database, or the data might be moved or flagged so that it could be recovered if the account is re-opened. A company might establish policies regarding salary maximums, minimums, and adjustments, in addition to basic sanity checks such as the salary not being a negative number. A retail database might not allow a purchase with the same serial number to be returned more than once, or might not allow credit card purchases above a certain value, while a database used to detect fraud might allow these kinds of things.
  See Also relational.

C

.cfg file
: A metadata file used with the `InnoDB` **transportable tablespace** feature. It is produced by the command `FLUSH TABLES ... FOR EXPORT`, puts one or more tables in a consistent state that can be copied to another server. The `.cfg` file is copied along with the corresponding **.ibd file**, and used to adjust the internal values of the `.ibd` file, such as the **space ID**, during the `ALTER TABLE ... IMPORT TABLESPACE` step.
  See Also .ibd file, space ID, transportable tablespace.

cache
: The general term for any memory area that stores copies of data for frequent or high-speed retrieval. In `InnoDB`, the primary kind of cache structure is the **buffer pool**.
  See Also buffer, buffer pool.

cardinality
: The number of different values in a table **column**. When queries refer to columns that have an associated **index**, the cardinality of each column influences which access method is most efficient. For example, for a column with a **unique constraint**, the number of different values is equal to the number of rows in the table. If a table has a million rows but only 10 different values for a particular column, each value occurs (on average) 100,000 times. A query such as `SELECT c1 FROM t1 WHERE c1 = 50;` thus might return 1 row or a huge number of rows, and the database server might process the query differently depending on the cardinality of `c1`.

  If the values in a column have a very uneven distribution, the cardinality might not be a good way to determine the best query plan. For example, `SELECT c1 FROM t1 WHERE c1 = x;` might return 1 row when `x=50` and a million rows when `x=30`. In such a case, you might need to use **index hints** to pass along advice about which lookup method is more efficient for a particular query.

  Cardinality can also apply to the number of distinct values present in multiple columns, as in a **composite index**.
  See Also column, composite index, index, index hint, persistent statistics, random dive, selectivity, unique constraint.

change buffer
: A special data structure that records changes to **pages** in **secondary indexes**. These values could result from SQL `INSERT`, `UPDATE`, or `DELETE` statements (**DML**). The set of features involving the change buffer is known collectively as **change buffering**, consisting of **insert buffering**, **delete buffering**, and **purge buffering**.

  Changes are only recorded in the change buffer when the relevant page from the secondary index is not in the **buffer pool**. When the relevant index page is brought into the buffer pool while associated changes are still in the change buffer, the changes for that page are applied in the buffer pool (**merged**) using the data

from the change buffer. Periodically, the **purge** operation that runs during times when the system is mostly idle, or during a slow shutdown, writes the new index pages to disk. The purge operation can write the disk blocks for a series of index values more efficiently than if each value were written to disk immediately.

Physically, the change buffer is part of the **system tablespace**, so that the index changes remain buffered across database restarts. The changes are only applied (**merged**) when the pages are brought into the buffer pool due to some other read operation.

The kinds and amount of data stored in the change buffer are governed by the `innodb_change_buffering` and `innodb_change_buffer_max_size` configuration options. To see information about the current data in the change buffer, issue the `SHOW ENGINE INNODB STATUS` command.

Formerly known as the **insert buffer**.
See Also buffer pool, change buffering, delete buffering, DML, insert buffer, insert buffering, merge, page, purge, purge buffering, secondary index, system tablespace.

change buffering
The general term for the features involving the **change buffer**, consisting of **insert buffering**, **delete buffering**, and **purge buffering**. Index changes resulting from SQL statements, which could normally involve random I/O operations, are held back and performed periodically by a background **thread**. This sequence of operations can write the disk blocks for a series of index values more efficiently than if each value were written to disk immediately. Controlled by the `innodb_change_buffering` and `innodb_change_buffer_max_size` configuration options.
See Also change buffer, delete buffering, insert buffering, purge buffering.

checkpoint
As changes are made to data pages that are cached in the **buffer pool**, those changes are written to the **data files** sometime later, a process known as **flushing**. The checkpoint is a record of the latest changes (represented by an **LSN** value) that have been successfully written to the data files.
See Also buffer pool, data files, flush, fuzzy checkpointing, LSN.

checksum
In `InnoDB`, a validation mechanism to detect corruption when a **page** in a **tablespace** is read from disk into the `InnoDB` **buffer pool**. This feature is controlled by the `innodb_checksums` configuration option in MySQL 5.5. `innodb_checksums` is deprecated in MySQL 5.6.3, replaced by `innodb_checksum_algorithm`.

The `innochecksum` command helps diagnose corruption problems by testing the checksum values for a specified **tablespace** file while the MySQL server is shut down.

MySQL also uses checksums for replication purposes. For details, see the configuration options `binlog_checksum`, `master_verify_checksum`, and `slave_sql_verify_checksum`.
See Also buffer pool, page, tablespace.

child table
In a **foreign key** relationship, a child table is one whose rows refer (or point) to rows in another table with an identical value for a specific column. This is the table that contains the `FOREIGN KEY ... REFERENCES` clause and optionally `ON UPDATE` and `ON DELETE` clauses. The corresponding row in the **parent table** must exist before the row can be created in the child table. The values in the child table can prevent delete or update operations on the parent table, or can cause automatic deletion or updates in the child table, based on the `ON CASCADE` option used when creating the foreign key.
See Also foreign key, parent table.

clean page
A **page** in the `InnoDB` **buffer pool** where all changes made in memory have also been written (**flushed**) to the data files. The opposite of a **dirty page**.
See Also buffer pool, data files, dirty page, flush, page.

clean shutdown
A **shutdown** that completes without errors and applies all changes to `InnoDB` tables before finishing, as opposed to a **crash** or a **fast shutdown**. Synonym for **slow shutdown**.

See Also crash, fast shutdown, shutdown, slow shutdown.

client
A type of program that sends requests to a **server**, and interprets or processes the results. The client software might run only some of the time (such as a mail or chat program), and might run interactively (such as the `mysql` command processor).
See Also mysql, server.

clustered index
The `InnoDB` term for a **primary key** index. `InnoDB` table storage is organized based on the values of the primary key columns, to speed up queries and sorts involving the primary key columns. For best performance, choose the primary key columns carefully based on the most performance-critical queries. Because modifying the columns of the clustered index is an expensive operation, choose primary columns that are rarely or never updated.

In the Oracle Database product, this type of table is known as an **index-organized table**.
See Also index, primary key, secondary index.

cold backup
A **backup** taken while the database is shut down. For busy applications and web sites, this might not be practical, and you might prefer a **warm backup** or a **hot backup**.
See Also backup, hot backup, warm backup.

column
A data item within a **row**, whose storage and semantics are defined by a data type. Each **table** and **index** is largely defined by the set of columns it contains.

Each column has a **cardinality** value. A column can be the **primary key** for its table, or part of the primary key. A column can be subject to a **unique constraint**, a **NOT NULL constraint**, or both. Values in different columns, even across different tables, can be linked by a **foreign key** relationship.

In discussions of MySQL internal operations, sometimes **field** is used as a synonym.
See Also cardinality, foreign key, index, NOT NULL constraint, primary key, row, table, unique constraint.

column index
An **index** on a single column.
See Also composite index, index.

column prefix
When an **index** is created with a length specification, such as `CREATE INDEX idx ON t1 (c1(N))`, only the first N characters of the column value are stored in the index. Keeping the index prefix small makes the index compact, and the memory and disk I/O savings help performance. (Although making the index prefix too small can hinder query optimization by making rows with different values appear to be duplicates.)

For columns containing binary values or long text strings, where sorting is not a major consideration and storing the entire value in the index would waste space, the index automatically uses the first N (typically 768) characters of the value to do lookups and sorts.
See Also index.

commit
A **SQL** statement that ends a **transaction**, making permanent any changes made by the transaction. It is the opposite of **rollback**, which undoes any changes made in the transaction.

`InnoDB` uses an **optimistic** mechanism for commits, so that changes can be written to the data files before the commit actually occurs. This technique makes the commit itself faster, with the tradeoff that more work is required in case of a rollback.

By default, MySQL uses the **autocommit** setting, which automatically issues a commit following each SQL statement.
See Also autocommit, optimistic, rollback, SQL, transaction.

compact row format

A **row format** for InnoDB tables. It was the default row format from MySQL 5.0.3 to MySQL 5.7.8. In MySQL 8.0, the default row format is defined by the `innodb_default_row_format` configuration option, which has a default setting of **DYNAMIC**. The **COMPACT** row format provides a more compact representation for nulls and variable-length columns than the **REDUNDANT** row format.

For additional information about `InnoDB` `COMPACT` row format, see Section 15.10.4, "COMPACT and REDUNDANT Row Formats".
See Also dynamic row format, file format, redundant row format, row format.

composite index

An **index** that includes multiple columns.
See Also index.

compressed backup

The compression feature of the **MySQL Enterprise Backup** product makes a compressed copy of each tablespace, changing the extension from `.ibd` to `.ibz`. Compressing backup data allows you to keep more backups on hand, and reduces the time to transfer backups to a different server. The data is uncompressed during the restore operation. When a compressed backup operation processes a table that is already compressed, it skips the compression step for that table, because compressing again would result in little or no space savings.

A set of files produced by the **MySQL Enterprise Backup** product, where each **tablespace** is compressed. The compressed files are renamed with a `.ibz` file extension.

Applying **compression** at the start of the backup process helps to avoid storage overhead during the compression process, and to avoid network overhead when transferring the backup files to another server. The process of **applying** the **binary log** takes longer, and requires uncompressing the backup files.
See Also apply, binary log, compression, hot backup, MySQL Enterprise Backup, tablespace.

compressed row format

A **row format** that enables data and index **compression** for `InnoDB` tables. Large fields are stored away from the page that holds the rest of the row data, as in **dynamic row format**. Both index pages and the large fields are compressed, yielding memory and disk savings. Depending on the structure of the data, the decrease in memory and disk usage might or might not outweigh the performance overhead of uncompressing the data as it is used. See Section 15.9, "InnoDB Table and Page Compression" for usage details.

For additional information about `InnoDB` `COMPRESSED` row format, see Section 15.10.3, "DYNAMIC and COMPRESSED Row Formats".
See Also compression, dynamic row format, row format.

compressed table

A table for which the data is stored in compressed form. For `InnoDB`, it is a table created with `ROW_FORMAT=COMPRESSED`. See Section 15.9, "InnoDB Table and Page Compression" for more information.
See Also compressed row format, compression.

compression

A feature with wide-ranging benefits from using less disk space, performing less I/O, and using less memory for caching.

`InnoDB` supports both table-level and page-level compression. `InnoDB` page compression is also referred to as **transparent page compression**. For more information about `InnoDB` compression, see Section 15.9, "InnoDB Table and Page Compression".

Another type of compression is the **compressed backup** feature of the **MySQL Enterprise Backup** product.
See Also buffer pool, compressed backup, compressed row format, DML, transparent page compression.

compression failure

Not actually an error, rather an expensive operation that can occur when using **compression** in combination with **DML** operations. It occurs when: updates to a compressed **page** overflow the area

on the page reserved for recording modifications; the page is compressed again, with all changes applied to the table data; the re-compressed data does not fit on the original page, requiring MySQL to split the data into two new pages and compress each one separately. To check the frequency of this condition, query the `INFORMATION_SCHEMA.INNODB_CMP` table and check how much the value of the `COMPRESS_OPS` column exceeds the value of the `COMPRESS_OPS_OK` column. Ideally, compression failures do not occur often; when they do, you can adjust the `innodb_compression_level`, `innodb_compression_failure_threshold_pct`, and `innodb_compression_pad_pct_max` configuration options.
See Also compression, DML, page.

concatenated index
See composite index.

concurrency
The ability of multiple operations (in database terminology, **transactions**) to run simultaneously, without interfering with each other. Concurrency is also involved with performance, because ideally the protection for multiple simultaneous transactions works with a minimum of performance overhead, using efficient mechanisms for **locking**.
See Also ACID, locking, transaction.

configuration file
The file that holds the **option** values used by MySQL at startup. Traditionally, on Linux and Unix this file is named `my.cnf`, and on Windows it is named `my.ini`. You can set a number of options related to InnoDB under the `[mysqld]` section of the file.

See Section 4.2.6, “Using Option Files” for information about where MySQL searches for configuration files.

When you use the **MySQL Enterprise Backup** product, you typically use two configuration files: one that specifies where the data comes from and how it is structured (which could be the original configuration file for your server), and a stripped-down one containing only a small set of options that specify where the backup data goes and how it is structured. The configuration files used with the **MySQL Enterprise Backup** product must contain certain options that are typically left out of regular configuration files, so you might need to add options to your existing configuration file for use with **MySQL Enterprise Backup**.
See Also my.cnf, MySQL Enterprise Backup, option, option file.

consistent read
A read operation that uses **snapshot** information to present query results based on a point in time, regardless of changes performed by other transactions running at the same time. If queried data has been changed by another transaction, the original data is reconstructed based on the contents of the **undo log**. This technique avoids some of the **locking** issues that can reduce **concurrency** by forcing transactions to wait for other transactions to finish.

With **REPEATABLE READ isolation level**, the snapshot is based on the time when the first read operation is performed. With **READ COMMITTED** isolation level, the snapshot is reset to the time of each consistent read operation.

Consistent read is the default mode in which `InnoDB` processes `SELECT` statements in **READ COMMITTED** and **REPEATABLE READ** isolation levels. Because a consistent read does not set any locks on the tables it accesses, other sessions are free to modify those tables while a consistent read is being performed on the table.

For technical details about the applicable isolation levels, see Section 15.5.2.3, “Consistent Nonlocking Reads”.
See Also concurrency, isolation level, locking, READ COMMITTED, REPEATABLE READ, snapshot, transaction, undo log.

constraint
An automatic test that can block database changes to prevent data from becoming inconsistent. (In computer science terms, a kind of assertion related to an invariant condition.) Constraints are a crucial component of

the **ACID** philosophy, to maintain data consistency. Constraints supported by MySQL include **FOREIGN KEY constraints** and **unique constraints**.
See Also ACID, foreign key, unique constraint.

counter
A value that is incremented by a particular kind of `InnoDB` operation. Useful for measuring how busy a server is, troubleshooting the sources of performance issues, and testing whether changes (for example, to configuration settings or indexes used by queries) have the desired low-level effects. Different kinds of counters are available through **Performance Schema** tables and **INFORMATION_SCHEMA** tables, particularly `INFORMATION_SCHEMA.INNODB_METRICS`.
See Also INFORMATION_SCHEMA, metrics counter, Performance Schema.

covering index
An **index** that includes all the columns retrieved by a query. Instead of using the index values as pointers to find the full table rows, the query returns values from the index structure, saving disk I/O. `InnoDB` can apply this optimization technique to more indexes than MyISAM can, because `InnoDB` **secondary indexes** also include the **primary key** columns. `InnoDB` cannot apply this technique for queries against tables modified by a transaction, until that transaction ends.

Any **column index** or **composite index** could act as a covering index, given the right query. Design your indexes and queries to take advantage of this optimization technique wherever possible.
See Also column index, composite index, index, primary key, secondary index.

CPU-bound
A type of **workload** where the primary **bottleneck** is CPU operations in memory. Typically involves read-intensive operations where the results can all be cached in the **buffer pool**.
See Also bottleneck, buffer pool, workload.

crash
MySQL uses the term “crash” to refer generally to any unexpected **shutdown** operation where the server cannot do its normal cleanup. For example, a crash could happen due to a hardware fault on the database server machine or storage device; a power failure; a potential data mismatch that causes the MySQL server to halt; a **fast shutdown** initiated by the DBA; or many other reasons. The robust, automatic **crash recovery** for **InnoDB** tables ensures that data is made consistent when the server is restarted, without any extra work for the DBA.
See Also crash recovery, fast shutdown, InnoDB, shutdown.

crash recovery
The cleanup activities that occur when MySQL is started again after a **crash**. For **InnoDB** tables, changes from incomplete transactions are replayed using data from the **redo log**. Changes that were **committed** before the crash, but not yet written into the **data files**, are reconstructed from the **doublewrite buffer**. When the database is shut down normally, this type of activity is performed during shutdown by the **purge** operation.

During normal operation, committed data can be stored in the **change buffer** for a period of time before being written to the data files. There is always a tradeoff between keeping the data files up-to-date, which introduces performance overhead during normal operation, and buffering the data, which can make shutdown and crash recovery take longer.
See Also change buffer, commit, crash, data files, doublewrite buffer, InnoDB, purge, redo log.

CRUD
Acronym for “create, read, update, delete”, a common sequence of operations in database applications. Often denotes a class of applications with relatively simple database usage (basic **DDL**, **DML** and **query** statements in **SQL**) that can be implemented quickly in any language.
See Also DDL, DML, query, SQL.

cursor
An internal data structure that is used to represent the result set of a **query**, or other operation that performs a search using an SQL `WHERE` clause. It works like an iterator in other high-level languages, producing each value from the result set as requested.

Although SQL usually handles the processing of cursors for you, you might delve into the inner workings when dealing with performance-critical code.
See Also query.

D

data definition language
See DDL.

data dictionary
Metadata that keeps track of database objects such as **tables**, **indexes**, and table **columns**. For the MySQL data dictionary, introduced in MySQL 8.0, metadata is physically located in `InnoDB` **file-per-table** tablespace files in the `mysql` database directory. For the `InnoDB` data dictionary, metadata is physically located in the `InnoDB` **system tablespace**.

Because the **MySQL Enterprise Backup** product always backs up the `InnoDB` system tablespace, all backups include the contents of the `InnoDB` data dictionary.
See Also column, file-per-table, .frm file, index, MySQL Enterprise Backup, system tablespace, table.

data directory
The directory under which each MySQL **instance** keeps the **data files** for `InnoDB` and the directories representing individual databases. Controlled by the `datadir` configuration option.
See Also data files, instance.

data files
The files that physically contain **table** and **index** data.

The `InnoDB` **system tablespace**, which holds the `InnoDB` **data dictionary** and is capable of holding data for multiple `InnoDB` tables, is represented by one or more `.ibdata` data files.

File-per-table tablespaces, which hold data for a single `InnoDB` table, are represented by a `.ibd` data file.

General tablespaces (introduced in MySQL 5.7.6), which can hold data for multiple `InnoDB` tables, are also represented by a `.ibd` data file.
See Also data dictionary, file-per-table, general tablespace, .ibd file, ibdata file, index, system tablespace, table, tablespace.

data manipulation language
See DML.

data warehouse
A database system or application that primarily runs large **queries**. The read-only or read-mostly data might be organized in **denormalized** form for query efficiency. Can benefit from the optimizations for **read-only transactions** in MySQL 5.6 and higher. Contrast with **OLTP**.
See Also denormalized, OLTP, query, read-only transaction.

database
Within the MySQL **data directory**, each database is represented by a separate directory. The InnoDB **system tablespace**, which can hold table data from multiple databases within a MySQL **instance**, is kept in **data files** that reside outside of individual database directories. When **file-per-table** mode is enabled, the **.ibd files** representing individual InnoDB tables are stored inside the database directories unless created elsewhere using the `DATA DIRECTORY` clause. General tablespaces, introduced in MySQL 5.7.6, also hold table data in **.ibd files**. Unlike file-per-table **.ibd files**, general tablespace **.ibd files** can hold table data from multiple databases within a MySQL **instance**, and can be assigned to directories relative to or independent of the MySQL data directory.

For long-time MySQL users, a database is a familiar notion. Users coming from an Oracle Database background will find that the MySQL meaning of a database is closer to what Oracle Database calls a **schema**.

See Also data files, file-per-table, .ibd file, instance, schema, system tablespace.

DCL
Data control language, a set of **SQL** statements for managing privileges. In MySQL, consists of the `GRANT` and `REVOKE` statements. Contrast with **DDL** and **DML**.
See Also DDL, DML, SQL.

DDL
Data definition language, a set of **SQL** statements for manipulating the database itself rather than individual table rows. Includes all forms of the `CREATE`, `ALTER`, and `DROP` statements. Also includes the `TRUNCATE` statement, because it works differently than a `DELETE FROM` *`table_name`* statement, even though the ultimate effect is similar.

DDL statements automatically **commit** the current **transaction**; they cannot be **rolled back**.

The `InnoDB` online DDL feature enhances performance for `CREATE INDEX`, `DROP INDEX`, and many types of `ALTER TABLE` operations. See Section 15.12, "InnoDB and Online DDL" for more information. Also, the `InnoDB` file-per-table setting can affect the behavior of `DROP TABLE` and `TRUNCATE TABLE` operations.

Contrast with **DML** and **DCL**.
See Also commit, DCL, DML, file-per-table, rollback, SQL, transaction.

deadlock
A situation where different **transactions** are unable to proceed, because each holds a **lock** that the other needs. Because both transactions are waiting for a resource to become available, neither will ever release the locks it holds.

A deadlock can occur when the transactions lock rows in multiple tables (through statements such as `UPDATE` or `SELECT ... FOR UPDATE`), but in the opposite order. A deadlock can also occur when such statements lock ranges of index records and **gaps**, with each transaction acquiring some locks but not others due to a timing issue.

For background information on how deadlocks are automatically detected and handled, see Section 15.5.5.2, "Deadlock Detection and Rollback". For tips on avoiding and recovering from deadlock conditions, see Section 15.5.5.3, "How to Minimize and Handle Deadlocks".
See Also gap, lock, transaction.

deadlock detection
A mechanism that automatically detects when a **deadlock** occurs, and automatically **rolls back** one of the **transactions** involved (the **victim**). Deadlock detection can be disabled using the `innodb_deadlock_detect` configuration option.
See Also deadlock, rollback, transaction, victim.

delete
When `InnoDB` processes a `DELETE` statement, the rows are immediately marked for deletion and no longer are returned by queries. The storage is reclaimed sometime later, during the periodic garbage collection known as the **purge** operation. For removing large quantities of data, related operations with their own performance characteristics are **TRUNCATE** and **DROP**.
See Also drop, purge, truncate.

delete buffering
The technique of storing changes to secondary index pages, resulting from `DELETE` operations, in the **change buffer** rather than writing the changes immediately, so that the physical writes can be performed to minimize random I/O. (Because delete operations are a two-step process, this operation buffers the write that normally marks an index record for deletion.) It is one of the types of **change buffering**; the others are **insert buffering** and **purge buffering**.
See Also change buffer, change buffering, insert buffer, insert buffering, purge buffering.

denormalized
A data storage strategy that duplicates data across different tables, rather than linking the tables with **foreign keys** and **join** queries. Typically used in **data warehouse** applications, where the data is not updated

after loading. In such applications, query performance is more important than making it simple to maintain consistent data during updates. Contrast with **normalized**.
See Also data warehouse, foreign key, join, normalized.

descending index
A type of **index** where index storage is optimized to process `ORDER BY column DESC` clauses.
See Also index.

dictionary object cache
The dictionary object cache stores previously accessed **data dictionary** objects in memory to enable object reuse and minimize disk I/O. An **LRU**-based eviction strategy is used to evict least recently used objects from memory. The cache is comprised of several partitions that store different object types.

For more information, see Section 14.4, “Dictionary Object Cache”.
See Also data dictionary, LRU.

dirty page
A **page** in the `InnoDB` **buffer pool** that has been updated in memory, where the changes are not yet written (**flushed**) to the **data files**. The opposite of a **clean page**.
See Also buffer pool, clean page, data files, flush, page.

dirty read
An operation that retrieves unreliable data, data that was updated by another transaction but not yet **committed**. It is only possible with the **isolation level** known as **read uncommitted**.

This kind of operation does not adhere to the **ACID** principle of database design. It is considered very risky, because the data could be **rolled back**, or updated further before being committed; then, the transaction doing the dirty read would be using data that was never confirmed as accurate.

Its opposite is **consistent read**, where `InnoDB` ensures that a transaction does not read information updated by another transaction, even if the other transaction commits in the meantime.
See Also ACID, commit, consistent read, isolation level, READ UNCOMMITTED, rollback.

disk-based
A kind of database that primarily organizes data on disk storage (hard drives or equivalent). Data is brought back and forth between disk and memory to be operated upon. It is the opposite of an **in-memory database**. Although `InnoDB` is disk-based, it also contains features such as he **buffer pool**, multiple buffer pool instances, and the **adaptive hash index** that allow certain kinds of workloads to work primarily from memory.
See Also adaptive hash index, buffer pool, in-memory database.

disk-bound
A type of **workload** where the primary **bottleneck** is disk I/O. (Also known as **I/O-bound**.) Typically involves frequent writes to disk, or random reads of more data than can fit into the **buffer pool**.
See Also bottleneck, buffer pool, workload.

DML
Data manipulation language, a set of **SQL** statements for performing `INSERT`, `UPDATE`, and `DELETE` operations. The `SELECT` statement is sometimes considered as a DML statement, because the `SELECT ... FOR UPDATE` form is subject to the same considerations for **locking** as `INSERT`, `UPDATE`, and `DELETE`.

DML statements for an `InnoDB` table operate in the context of a **transaction**, so their effects can be **committed** or **rolled back** as a single unit.

Contrast with **DDL** and **DCL**.
See Also commit, DCL, DDL, locking, rollback, SQL, transaction.

document id
In the `InnoDB` **full-text search** feature, a special column in the table containing the **FULLTEXT index**, to uniquely identify the document associated with each **ilist** value. Its name is `FTS_DOC_ID` (uppercase

required). The column itself must be of `BIGINT UNSIGNED NOT NULL` type, with a unique index named `FTS_DOC_ID_INDEX`. Preferably, you define this column when creating the table. If `InnoDB` must add the column to the table while creating a `FULLTEXT` index, the indexing operation is considerably more expensive.
See Also full-text search, FULLTEXT index, ilist.

doublewrite buffer
`InnoDB` uses a file flush technique called doublewrite. Before writing **pages** to the **data files**, `InnoDB` first writes them to a contiguous area called the doublewrite buffer. Only after the write and the flush to the doublewrite buffer have completed, does `InnoDB` write the pages to their proper positions in the data file. If there is an operating system, storage subsystem, or `mysqld` process crash in the middle of a page write, `InnoDB` can later find a good copy of the page from the doublewrite buffer during **crash recovery**.

Although data is always written twice, the doublewrite buffer does not require twice as much I/O overhead or twice as many I/O operations. Data is written to the buffer itself as a large sequential chunk, with a single `fsync()` call to the operating system.

To turn off the doublewrite buffer, specify the option `innodb_doublewrite=0`.
See Also crash recovery, data files, page, purge.

drop
A kind of **DDL** operation that removes a schema object, through a statement such as `DROP TABLE` or `DROP INDEX`. It maps internally to an `ALTER TABLE` statement. From an `InnoDB` perspective, the performance considerations of such operations involve the time that the **data dictionary** is locked to ensure that interrelated objects are all updated, and the time to update memory structures such as the **buffer pool**. For a **table**, the drop operation has somewhat different characteristics than a **truncate** operation (`TRUNCATE TABLE` statement).
See Also buffer pool, data dictionary, DDL, table, truncate.

dynamic row format
An `InnoDB` row format. Because long variable-length column values are stored outside of the page that holds the row data, it is very efficient for rows that include large objects. Since the large fields are typically not accessed to evaluate query conditions, they are not brought into the **buffer pool** as often, resulting in fewer I/O operations and better utilization of cache memory.

As of MySQL 5.7.9, the default row format is defined by `innodb_default_row_format`, which has a default value of `DYNAMIC`.

For additional information about `InnoDB` `DYNAMIC` row format, see Section 15.10.3, “DYNAMIC and COMPRESSED Row Formats”.
See Also buffer pool, file format, row format.

E

early adopter
A stage similar to **beta**, when a software product is typically evaluated for performance, functionality, and compatibility in a non-mission-critical setting.
See Also beta.

error log
A type of **log** showing information about MySQL startup and critical runtime errors and **crash** information. For details, see Section 5.4.2, “The Error Log”.
See Also crash, log.

eviction
The process of removing an item from a cache or other temporary storage area, such as the `InnoDB` **buffer pool**. Often, but not always, uses the **LRU** algorithm to determine which item to remove. When a **dirty page** is evicted, its contents are **flushed** to disk, and any dirty **neighbor pages** might be flushed also.
See Also buffer pool, dirty page, flush, LRU, neighbor page.

exclusive lock
A kind of **lock** that prevents any other **transaction** from locking the same row. Depending on the transaction **isolation level**, this kind of lock might block other transactions from writing to the same row, or might also block other transactions from reading the same row. The default `InnoDB` isolation level, **REPEATABLE READ**, enables higher **concurrency** by allowing transactions to read rows that have exclusive locks, a technique known as **consistent read**.
See Also concurrency, consistent read, isolation level, lock, REPEATABLE READ, shared lock, transaction.

extent
A group of **pages** within a **tablespace**. For the default **page size** of 16KB, an extent contains 64 pages. In MySQL 5.6, the page size for an `InnoDB` instance can be 4KB, 8KB, or 16KB, controlled by the `innodb_page_size` configuration option. For 4KB, 8KB, and 16KB pages sizes, the extent size is always 1MB (or 1048576 bytes).

Support for 32KB and 64KB `InnoDB` page sizes was added in MySQL 5.7.6. For a 32KB page size, the extent size is 2MB. For a 64KB page size, the extent size is 4MB.

`InnoDB` features such as **segments**, **read-ahead** requests and the **doublewrite buffer** use I/O operations that read, write, allocate, or free data one extent at a time.
See Also doublewrite buffer, page, page size, read-ahead, segment, tablespace.

F

.frm file
A file containing the metadata, such as the table definition, of a MySQL table. `.frm` files were removed in MySQL 8.0 but are still used in earlier MySQL releases. In MySQL 8.0, data previously stored in `.frm` files is stored in **data dictionary** tables.
See Also data dictionary, MySQL Enterprise Backup, system tablespace.

fanout
The increase or decrease in number of rows from adding a table during join processing. Suppose that a query of the form `SELECT ... FROM t1 INNER JOIN t2 ...` selects rows from table `t1`, then `t2`.

- If the join selects one row from `t2` per row in `t1`, the fanout for `t2` is 1.
- If the join selects two rows from `t2` per row in `t1`, the fanout for `t2` is 2.
- If the join selects one row from `t2` per two rows in `t1`, the fanout for `t2` is .5.

Fast Index Creation
A capability first introduced in the InnoDB Plugin, now part of MySQL in 5.5 and higher, that speeds up creation of `InnoDB` **secondary indexes** by avoiding the need to completely rewrite the associated table. The speedup applies to dropping secondary indexes also.

Because index maintenance can add performance overhead to many data transfer operations, consider doing operations such as `ALTER TABLE ... ENGINE=INNODB` or `INSERT INTO ... SELECT * FROM ...` without any secondary indexes in place, and creating the indexes afterward.

In MySQL 5.6, this feature becomes more general. You can read and write to tables while an index is being created, and many more kinds of `ALTER TABLE` operations can be performed without copying the table, without blocking **DML** operations, or both. Thus in MySQL 5.6 and higher, this set of features is referred to as **online DDL** rather than Fast Index Creation.

For related information, see InnoDB Fast Index Creation, and Section 15.12, “InnoDB and Online DDL”.
See Also DML, index, online DDL, secondary index.

fast shutdown
The default **shutdown** procedure for `InnoDB`, based on the configuration setting `innodb_fast_shutdown=1`. To save time, certain **flush** operations are skipped. This type of shutdown

is safe during normal usage, because the flush operations are performed during the next startup, using the same mechanism as in **crash recovery**. In cases where the database is being shut down for an upgrade or downgrade, do a **slow shutdown** instead to ensure that all relevant changes are applied to the **data files** during the shutdown.
See Also crash recovery, data files, flush, shutdown, slow shutdown.

file format
The file format for `InnoDB` tables.
See Also file-per-table, .ibd file, ibdata file, row format.

file-per-table
A general name for the setting controlled by the `innodb_file_per_table` option, which is an important configuration option that affects aspects of `InnoDB` file storage, availability of features, and I/O characteristics. As of MySQL 5.6.7, `innodb_file_per_table` is enabled by default.

With the `innodb_file_per_table` option enabled, you can create a table in its own **.ibd file** rather than in the shared **ibdata files** of the **system tablespace**. When table data is stored in an individual **.ibd file**, you have more flexibility to choose **row formats** required for features such as data **compression**. The `TRUNCATE TABLE` operation is also faster, and reclaimed space can be used by the operating system rather than remaining reserved for `InnoDB`.

The **MySQL Enterprise Backup** product is more flexible for tables that are in their own files. For example, tables can be excluded from a backup, but only if they are in separate files. Thus, this setting is suitable for tables that are backed up less frequently or on a different schedule.
See Also compressed row format, compression, file format, .ibd file, ibdata file, innodb_file_per_table, MySQL Enterprise Backup, row format, system tablespace.

fill factor
In an `InnoDB` **index**, the proportion of a **page** that is taken up by index data before the page is split. The unused space when index data is first divided between pages allows for rows to be updated with longer string values without requiring expensive index maintenance operations. If the fill factor is too low, the index consumes more space than needed, causing extra I/O overhead when reading the index. If the fill factor is too high, any update that increases the length of column values can cause extra I/O overhead for index maintenance. See Section 15.8.2.2, “The Physical Structure of an InnoDB Index” for more information.
See Also index, page.

fixed row format
This row format is used by the `MyISAM` storage engine, not by `InnoDB`. If you create an `InnoDB` table with the option `ROW_FORMAT=FIXED` in MySQL 5.7.6 or earlier, `InnoDB` uses the **compact row format** instead, although the `FIXED` value might still show up in output such as `SHOW TABLE STATUS` reports. As of MySQL 5.7.7, `InnoDB` returns an error if `ROW_FORMAT=FIXED` is specified.
See Also compact row format, row format.

flush
To write changes to the database files, that had been buffered in a memory area or a temporary disk storage area. The `InnoDB` storage structures that are periodically flushed include the **redo log**, the **undo log**, and the **buffer pool**.

Flushing can happen because a memory area becomes full and the system needs to free some space, because a **commit** operation means the changes from a transaction can be finalized, or because a **slow shutdown** operation means that all outstanding work should be finalized. When it is not critical to flush all the buffered data at once, `InnoDB` can use a technique called **fuzzy checkpointing** to flush small batches of pages to spread out the I/O overhead.
See Also buffer pool, commit, fuzzy checkpointing, redo log, slow shutdown, undo log.

flush list
An internal `InnoDB` data structure that tracks **dirty pages** in the **buffer pool**: that is, **pages** that have been changed and need to be written back out to disk. This data structure is updated frequently by `InnoDB` internal **mini-transactions**, and so is protected by its own **mutex** to allow concurrent access to the buffer pool.

See Also buffer pool, dirty page, LRU, mini-transaction, mutex, page, page cleaner.

foreign key
A type of pointer relationship, between rows in separate `InnoDB` tables. The foreign key relationship is defined on one column in both the **parent table** and the **child table**.

In addition to enabling fast lookup of related information, foreign keys help to enforce **referential integrity**, by preventing any of these pointers from becoming invalid as data is inserted, updated, and deleted. This enforcement mechanism is a type of **constraint**. A row that points to another table cannot be inserted if the associated foreign key value does not exist in the other table. If a row is deleted or its foreign key value changed, and rows in another table point to that foreign key value, the foreign key can be set up to prevent the deletion, cause the corresponding column values in the other table to become **null**, or automatically delete the corresponding rows in the other table.

One of the stages in designing a **normalized** database is to identify data that is duplicated, separate that data into a new table, and set up a foreign key relationship so that the multiple tables can be queried like a single table, using a **join** operation.
See Also child table, FOREIGN KEY constraint, join, normalized, NULL, parent table, referential integrity, relational.

FOREIGN KEY constraint
The type of **constraint** that maintains database consistency through a **foreign key** relationship. Like other kinds of constraints, it can prevent data from being inserted or updated if data would become inconsistent; in this case, the inconsistency being prevented is between data in multiple tables. Alternatively, when a **DML** operation is performed, `FOREIGN KEY` constraints can cause data in **child rows** to be deleted, changed to different values, or set to **null**, based on the `ON CASCADE` option specified when creating the foreign key.
See Also child table, constraint, DML, foreign key, NULL.

FTS
In most contexts, an acronym for **full-text search**. Sometimes in performance discussions, an acronym for **full table scan**.
See Also full table scan, full-text search.

full backup
A **backup** that includes all the **tables** in each MySQL **database**, and all the databases in a MySQL **instance**. Contrast with **partial backup**.
See Also backup, database, instance, partial backup, table.

full table scan
An operation that requires reading the entire contents of a table, rather than just selected portions using an **index**. Typically performed either with small lookup tables, or in data warehousing situations with large tables where all available data is aggregated and analyzed. How frequently these operations occur, and the sizes of the tables relative to available memory, have implications for the algorithms used in query optimization and managing the **buffer pool**.

The purpose of indexes is to allow lookups for specific values or ranges of values within a large table, thus avoiding full table scans when practical.
See Also buffer pool, index.

full-text search
The MySQL feature for finding words, phrases, Boolean combinations of words, and so on within table data, in a faster, more convenient, and more flexible way than using the SQL `LIKE` operator or writing your own application-level search algorithm. It uses the SQL function `MATCH()` and **FULLTEXT indexes**.
See Also FULLTEXT index.

FULLTEXT index
The special kind of **index** that holds the **search index** in the MySQL **full-text search** mechanism. Represents the words from values of a column, omitting any that are specified as **stopwords**. Originally, only available for `MyISAM` tables. Starting in MySQL 5.6.4, it is also available for **InnoDB** tables.

See Also full-text search, index, InnoDB, search index, stopword.

fuzzy checkpointing
A technique that **flushes** small batches of **dirty pages** from the **buffer pool**, rather than flushing all dirty pages at once which would disrupt database processing.
See Also buffer pool, dirty page, flush.

G

GA
"Generally available", the stage when a software product leaves **beta** and is available for sale, official support, and production use.
See Also beta.

gap
A place in an `InnoDB` **index** data structure where new values could be inserted. When you lock a set of rows with a statement such as `SELECT ... FOR UPDATE`, `InnoDB` can create locks that apply to the gaps as well as the actual values in the index. For example, if you select all values greater than 10 for update, a gap lock prevents another transaction from inserting a new value that is greater than 10. The **supremum record** and **infimum record** represent the gaps containing all values greater than or less than all the current index values.
See Also concurrency, gap lock, index, infimum record, isolation level, supremum record.

gap lock
A **lock** on a **gap** between index records, or a lock on the gap before the first or after the last index record. For example, `SELECT c1 FROM t WHERE c1 BETWEEN 10 and 20 FOR UPDATE;` prevents other transactions from inserting a value of 15 into the column `t.c1`, whether or not there was already any such value in the column, because the gaps between all existing values in the range are locked. Contrast with **record lock** and **next-key lock**.

Gap locks are part of the tradeoff between performance and **concurrency**, and are used in some transaction **isolation levels** and not others.
See Also gap, infimum record, lock, next-key lock, record lock, supremum record.

general log
See general query log.

general query log
A type of **log** used for diagnosis and troubleshooting of SQL statements processed by the MySQL server. Can be stored in a file or in a database table. You must enable this feature through the `general_log` configuration option to use it. You can disable it for a specific connection through the `sql_log_off` configuration option.

Records a broader range of queries than the **slow query log**. Unlike the **binary log**, which is used for replication, the general query log contains `SELECT` statements and does not maintain strict ordering. For more information, see Section 5.4.3, "The General Query Log".
See Also binary log, log, slow query log.

general tablespace
A shared `InnoDB` **tablespace** created using `CREATE TABLESPACE` syntax. General tablespaces can be created outside of the MySQL data directory, are capable of holding multiple **tables**, and support tables of all row formats. General tablespaces were introduced in MySQL 5.7.6.

Tables are added to a general tablespace using `CREATE TABLE` *`tbl_name`* `... TABLESPACE [=]` *`tablespace_name`* or `ALTER TABLE` *`tbl_name`* `TABLESPACE [=]` *`tablespace_name`* syntax.

Contrast with **system tablespace** and **file-per-table** tablespace.

For more information, see Section 15.7.9, "InnoDB General Tablespaces".
See Also file-per-table, system tablespace, table, tablespace.

generated column
A column whose values are computed from an expression included in the column definition. A generated column can be **virtual** or **stored**.
See Also base column, stored generated column, virtual generated column.

generated stored column
See stored generated column.

generated virtual column
See virtual generated column.

global transaction
A type of **transaction** involved in **XA** operations. It consists of several actions that are transactional in themselves, but that all must either complete successfully as a group, or all be rolled back as a group. In essence, this extends **ACID** properties "up a level" so that multiple ACID transactions can be executed in concert as components of a global operation that also has ACID properties.
See Also ACID, transaction, XA.

group commit
An `InnoDB` optimization that performs some low-level I/O operations (log write) once for a set of **commit** operations, rather than flushing and syncing separately for each commit.

When the **binary log** is enabled, you typically also set the configuration option `sync_binlog=0`, because group commit for the binary log is only supported if it is set to 0.
See Also binary log, commit.

H

hash index
A type of **index** intended for queries that use equality operators, rather than range operators such as greater-than or `BETWEEN`. It is available for `MEMORY` tables. Although hash indexes are the default for `MEMORY` tables for historic reasons, that storage engine also supports **B-tree** indexes, which are often a better choice for general-purpose queries.

MySQL includes a variant of this index type, the **adaptive hash index**, that is constructed automatically for `InnoDB` tables if needed based on runtime conditions.
See Also adaptive hash index, B-tree, index, InnoDB.

HDD
Acronym for "hard disk drive". Refers to storage media using spinning platters, usually when comparing and contrasting with **SSD**. Its performance characteristics can influence the throughput of a **disk-based** workload.
See Also disk-based, SSD.

heartbeat
A periodic message that is sent to indicate that a system is functioning properly. In a **replication** context, if the **master** stops sending such messages, one of the **slaves** can take its place. Similar techniques can be used between the servers in a cluster environment, to confirm that all of them are operating properly.
See Also master server, replication.

high-water mark
A value representing an upper limit, either a hard limit that should not be exceeded at runtime, or a record of the maximum value that was actually reached. Contrast with **low-water mark**.
See Also low-water mark.

history list
A list of **transactions** with delete-marked records scheduled to be processed by the `InnoDB` **purge** operation. Recorded in the **undo log**. The length of the history list is reported by the command `SHOW ENGINE INNODB STATUS`. If the history list grows longer than the value of the `innodb_max_purge_lag` configuration option, each **DML** operation is delayed slightly to allow the purge operation to finish **flushing** the deleted records.

Also known as **purge lag**.
See Also DML, flush, purge, purge lag, rollback segment, transaction, undo log.

hole punching
Releasing empty blocks from a page. The `InnoDB` **transparent page compression** feature relies on hole punching support. For more information, see Section 15.9.2, “InnoDB Page Compression”.
See Also sparse file, transparent page compression.

hot
A condition where a row, table, or internal data structure is accessed so frequently, requiring some form of locking or mutual exclusion, that it results in a performance or scalability issue.

Although “hot” typically indicates an undesirable condition, a **hot backup** is the preferred type of backup.
See Also hot backup.

hot backup
A backup taken while the database is running and applications are reading and writing to it. The backup involves more than simply copying data files: it must include any data that was inserted or updated while the backup was in process; it must exclude any data that was deleted while the backup was in process; and it must ignore any changes that were not committed.

The Oracle product that performs hot backups, of `InnoDB` tables especially but also tables from `MyISAM` and other storage engines, is known as **MySQL Enterprise Backup**.

The hot backup process consists of two stages. The initial copying of the data files produces a **raw backup**. The **apply** step incorporates any changes to the database that happened while the backup was running. Applying the changes produces a **prepared** backup; these files are ready to be restored whenever necessary.
See Also apply, MySQL Enterprise Backup, prepared backup, raw backup.

I

.ibd file
The data file for **file-per-table** tablespaces and general tablespaces. File-per-table tablespace `.idb` files contain a single table and associated index data. **General tablespace** `.idb` files may contain table and index data for multiple tables.

The `.ibd` file extension does not apply to the **system tablespace**, which consists of one or more **ibdata files**.

If a file-per-table tablespace or general tablespace is created with the `DATA DIRECTORY =` clause, the `.ibd` file is located at the specified path, outside the normal data directory.

When a `.ibd` file is included in a compressed backup by the **MySQL Enterprise Backup** product, the compressed equivalent is a `.ibz` file.
See Also database, file-per-table, general tablespace, ibdata file, .ibz file, innodb_file_per_table, MySQL Enterprise Backup, system tablespace.

.ibz file
When the **MySQL Enterprise Backup** product performs a **compressed backup**, it transforms each **tablespace** file that is created using the **file-per-table** setting from a `.ibd` extension to a `.ibz` extension.

The compression applied during backup is distinct from the **compressed row format** that keeps table data compressed during normal operation. A compressed backup operation skips the compression step for a tablespace that is already in compressed row format, as compressing a second time would slow down the backup but produce little or no space savings.
See Also compressed backup, compressed row format, file-per-table, .ibd file, MySQL Enterprise Backup, tablespace.

I/O-bound
See disk-bound.

ib-file set

The set of files managed by InnoDB within a MySQL database: the **system tablespace**, **file-per-table** tablespace files, and **redo log** files. Depending on MySQL version and InnoDB configuration, may also include **general tablespace**, **temporary tablespace**, and **undo tablespace** files. This term is sometimes used in detailed discussions of InnoDB file structures and formats to refer to the set of files managed by InnoDB within a MySQL database.

See Also database, file-per-table, general tablespace, redo log, system tablespace, temporary tablespace, undo tablespace.

ibbackup_logfile

A supplemental backup file created by the **MySQL Enterprise Backup** product during a **hot backup** operation. It contains information about any data changes that occurred while the backup was running. The initial backup files, including ibbackup_logfile, are known as a **raw backup**, because the changes that occurred during the backup operation are not yet incorporated. After you perform the **apply** step to the raw backup files, the resulting files do include those final data changes, and are known as a **prepared backup**. At this stage, the ibbackup_logfile file is no longer necessary.

See Also apply, hot backup, MySQL Enterprise Backup, prepared backup, raw backup.

ibdata file

A set of files with names such as ibdata1, ibdata2, and so on, that make up the InnoDB **system tablespace**. These files contain metadata about InnoDB tables, (the InnoDB **data dictionary**), and the storage areas for one or more **undo logs**, the **change buffer**, and the **doublewrite buffer**. They also can contain some or all of the table data also (depending on whether the **file-per-table** mode is in effect when each table is created). When the **innodb_file_per_table** option is enabled, data and indexes for newly created tables are stored in separate **.ibd files** rather than in the system tablespace.

The growth of the ibdata files is influenced by the innodb_autoextend_increment configuration option.

See Also change buffer, data dictionary, doublewrite buffer, file-per-table, .ibd file, innodb_file_per_table, system tablespace, undo log.

ibtmp file

The InnoDB **temporary tablespace data file** for non-compressed InnoDB **temporary tables** and related objects. The configuration file option, innodb_temp_data_file_path, allows users to define a relative path for the temporary tablespace data file. If innodb_temp_data_file_path is not specified, the default behavior is to create a single auto-extending 12MB data file named ibtmp1 in the data directory, alongside ibdata1.

See Also data files, temporary table, temporary tablespace.

ib_logfile

A set of files, typically named ib_logfile0 and ib_logfile1, that form the **redo log**. Also sometimes referred to as the **log group**. These files record statements that attempt to change data in InnoDB tables. These statements are replayed automatically to correct data written by incomplete transactions, on startup following a crash.

This data cannot be used for manual recovery; for that type of operation, use the **binary log**.

See Also binary log, log group, redo log.

ilist

Within an InnoDB **FULLTEXT index**, the data structure consisting of a document ID and positional information for a token (that is, a particular word).

See Also FULLTEXT index.

implicit row lock

A row lock that InnoDB acquires to ensure consistency, without you specifically requesting it.

See Also row lock.

in-memory database

A type of database system that maintains data in memory, to avoid overhead due to disk I/O and translation between disk blocks and memory areas. Some in-memory databases sacrifice durability (the “D” in the **ACID**

design philosophy) and are vulnerable to hardware, power, and other types of failures, making them more suitable for read-only operations. Other in-memory databases do use durability mechanisms such as logging changes to disk or using non-volatile memory.

MySQL features that address the same kinds of memory-intensive processing include the `InnoDB` **buffer pool**, **adaptive hash index**, and **read-only transaction** optimization, the `MEMORY` storage engine, the `MyISAM` key cache, and the MySQL query cache.
See Also ACID, adaptive hash index, buffer pool, disk-based, read-only transaction.

incremental backup
A type of **hot backup**, performed by the **MySQL Enterprise Backup** product, that only saves data changed since some point in time. Having a full backup and a succession of incremental backups lets you reconstruct backup data over a long period, without the storage overhead of keeping several full backups on hand. You can restore the full backup and then apply each of the incremental backups in succession, or you can keep the full backup up-to-date by applying each incremental backup to it, then perform a single restore operation.

The granularity of changed data is at the **page** level. A page might actually cover more than one row. Each changed page is included in the backup.
See Also hot backup, MySQL Enterprise Backup, page.

index
A data structure that provides a fast lookup capability for **rows** of a **table**, typically by forming a tree structure (**B-tree)** representing all the values of a particular **column** or set of columns.

`InnoDB` tables always have a **clustered index** representing the **primary key**. They can also have one or more **secondary indexes** defined on one or more columns. Depending on their structure, secondary indexes can be classified as **partial**, **column**, or **composite** indexes.

Indexes are a crucial aspect of **query** performance. Database architects design tables, queries, and indexes to allow fast lookups for data needed by applications. The ideal database design uses a **covering index** where practical; the query results are computed entirely from the index, without reading the actual table data. Each **foreign key** constraint also requires an index, to efficiently check whether values exist in both the **parent** and **child** tables.

Although a B-tree index is the most common, a different kind of data structure is used for **hash indexes**, as in the `MEMORY` storage engine and the `InnoDB` **adaptive hash index**. **R-tree** indexes are used for spatial indexing of multi-dimensional information.
See Also adaptive hash index, B-tree, child table, clustered index, column index, composite index, covering index, foreign key, hash index, parent table, partial index, primary key, query, R-tree, row, secondary index, table.

index cache
A memory area that holds the token data for `InnoDB` **full-text search**. It buffers the data to minimize disk I/O when data is inserted or updated in columns that are part of a **FULLTEXT index**. The token data is written to disk when the index cache becomes full. Each `InnoDB` `FULLTEXT` index has its own separate index cache, whose size is controlled by the configuration option `innodb_ft_cache_size`.
See Also full-text search, FULLTEXT index.

index condition pushdown
Index condition pushdown (ICP) is an optimization that pushes part of a `WHERE` condition down to the storage engine if parts of the condition can be evaluated using fields from the **index**. ICP can reduce the number of times the **storage engine** must access the base table and the number of times the MySQL server must access the storage engine. For more information, see Section 8.2.1.5, “Index Condition Pushdown Optimization”.
See Also index, storage engine.

index hint
Extended SQL syntax for overriding the **indexes** recommended by the optimizer. For example, the `FORCE INDEX`, `USE INDEX`, and `IGNORE INDEX` clauses. Typically used when indexed columns have unevenly distributed values, resulting in inaccurate **cardinality** estimates.

See Also cardinality, index.

index prefix
In an **index** that applies to multiple columns (known as a **composite index**), the initial or leading columns of the index. A query that references the first 1, 2, 3, and so on columns of a composite index can use the index, even if the query does not reference all the columns in the index.
See Also composite index, index.

index statistics
See statistics.

infimum record
A **pseudo-record** in an **index**, representing the **gap** below the smallest value in that index. If a transaction has a statement such as `SELECT ... FROM ... WHERE col < 10 FOR UPDATE;`, and the smallest value in the column is 5, it is a lock on the infimum record that prevents other transactions from inserting even smaller values such as 0, -10, and so on.
See Also gap, index, pseudo-record, supremum record.

INFORMATION_SCHEMA
The name of the **database** that provides a query interface to the MySQL **data dictionary**. (This name is defined by the ANSI SQL standard.) To examine information (metadata) about the database, you can query tables such as `INFORMATION_SCHEMA.TABLES` and `INFORMATION_SCHEMA.COLUMNS`, rather than using `SHOW` commands that produce unstructured output.

The `INFORMATION_SCHEMA` database also contains tables specific to **InnoDB** that provide a query interface to the `InnoDB` data dictionary. You use these tables not to see how the database is structured, but to get real-time information about the workings of `InnoDB` tables to help with performance monitoring, tuning, and troubleshooting.
See Also data dictionary, database, InnoDB.

InnoDB
A MySQL component that combines high performance with **transactional** capability for reliability, robustness, and concurrent access. It embodies the **ACID** design philosophy. Represented as a **storage engine**; it handles tables created or altered with the `ENGINE=INNODB` clause. See Chapter 15, *The InnoDB Storage Engine* for architectural details and administration procedures, and Section 8.5, “Optimizing for InnoDB Tables” for performance advice.

In MySQL 5.5 and higher, `InnoDB` is the default storage engine for new tables and the `ENGINE=INNODB` clause is not required.

`InnoDB` tables are ideally suited for **hot backups**. See Section 29.2, “MySQL Enterprise Backup Overview” for information about the **MySQL Enterprise Backup** product for backing up MySQL servers without interrupting normal processing.
See Also ACID, hot backup, MySQL Enterprise Backup, storage engine, transaction.

innodb_autoinc_lock_mode
The `innodb_autoinc_lock_mode` option controls the algorithm used for **auto-increment locking**. When you have an auto-incrementing **primary key**, you can use statement-based replication only with the setting `innodb_autoinc_lock_mode=1`. This setting is known as *consecutive* lock mode, because multi-row inserts within a transaction receive consecutive auto-increment values. If you have `innodb_autoinc_lock_mode=2`, which allows higher concurrency for insert operations, use row-based replication rather than statement-based replication. This setting is known as *interleaved* lock mode, because multiple multi-row insert statements running at the same time can receive **auto-increment** values that are interleaved. The setting `innodb_autoinc_lock_mode=0` is the previous (traditional) default setting and should not be used except for compatibility purposes.
See Also auto-increment, auto-increment locking, mixed-mode insert, primary key.

innodb_file_per_table
An important configuration option that affects many aspects of `InnoDB` file storage, availability of features, and I/O characteristics. In MySQL 5.6.7 and higher, it is enabled by default. The `innodb_file_per_table`

option turns on **file-per-table** mode. With this mode enabled, a newly created `InnoDB` table and associated indexes can be stored in a file-per-table **.ibd file**, outside the **system tablespace**.

This option affects the performance and storage considerations for a number of SQL statements, such as `DROP TABLE` and `TRUNCATE TABLE`.

Enabling the `innodb_file_per_table` option allows you to take advantage of features such as table **compression** and named-table backups in **MySQL Enterprise Backup**.

For more information, see `innodb_file_per_table`, and Section 15.7.4, “InnoDB File-Per-Table Tablespaces”.
See Also compression, file-per-table, .ibd file, MySQL Enterprise Backup, system tablespace.

innodb_lock_wait_timeout
The `innodb_lock_wait_timeout` option sets the balance between **waiting** for shared resources to become available, or giving up and handling the error, retrying, or doing alternative processing in your application. Rolls back any `InnoDB` transaction that waits more than a specified time to acquire a **lock**. Especially useful if **deadlocks** are caused by updates to multiple tables controlled by different storage engines; such deadlocks are not **detected** automatically.
See Also deadlock, deadlock detection, lock, wait.

innodb_strict_mode
The `innodb_strict_mode` option controls whether `InnoDB` operates in **strict mode**, where conditions that are normally treated as warnings, cause errors instead (and the underlying statements fail).
See Also strict mode.

insert
One of the primary **DML** operations in **SQL**. The performance of inserts is a key factor in **data warehouse** systems that load millions of rows into tables, and **OLTP** systems where many concurrent connections might insert rows into the same table, in arbitrary order. If insert performance is important to you, you should learn about **InnoDB** features such as the **insert buffer** used in **change buffering**, and **auto-increment** columns.
See Also auto-increment, change buffering, data warehouse, DML, InnoDB, insert buffer, OLTP, SQL.

insert buffer
The former name of the **change buffer**. In MySQL 5.5, support was added for buffering changes to secondary index pages for `DELETE` and `UPDATE` operations. Previously, only changes resulting from `INSERT` operations were buffered. The preferred term is now *change buffer*.
See Also change buffer, change buffering.

insert buffering
The technique of storing changes to secondary index pages, resulting from `INSERT` operations, in the **change buffer** rather than writing the changes immediately, so that the physical writes can be performed to minimize random I/O. It is one of the types of **change buffering**; the others are **delete buffering** and **purge buffering**.

Insert buffering is not used if the secondary index is **unique**, because the uniqueness of new values cannot be verified before the new entries are written out. Other kinds of change buffering do work for unique indexes.
See Also change buffer, change buffering, delete buffering, insert buffer, purge buffering, unique index.

insert intention lock
A type of **gap lock** that is set by `INSERT` operations prior to row insertion. This type of **lock** signals the intent to insert in such a way that multiple transactions inserting into the same index gap need not wait for each other if they are not inserting at the same position within the gap. For more information, see Section 15.5.1, “InnoDB Locking”.
See Also gap lock, lock, next-key lock.

instance
A single **mysqld** daemon managing a **data directory** representing one or more **databases** with a set of **tables**. It is common in development, testing, and some **replication** scenarios to have multiple instances on the same **server** machine, each managing its own data directory and listening on its own port or socket. With

one instance running a **disk-bound** workload, the server might still have extra CPU and memory capacity to run additional instances.
See Also data directory, database, disk-bound, mysqld, replication, server, table.

instrumentation
Modifications at the source code level to collect performance data for tuning and debugging. In MySQL, data collected by instrumentation is exposed through an SQL interface using the `INFORMATION_SCHEMA` and `PERFORMANCE_SCHEMA` databases.
See Also INFORMATION_SCHEMA, Performance Schema.

intention exclusive lock
See intention lock.

intention lock
A kind of **lock** that applies to the table, used to indicate the kind of lock the **transaction** intends to acquire on rows in the table. Different transactions can acquire different kinds of intention locks on the same table, but the first transaction to acquire an *intention exclusive* (IX) lock on a table prevents other transactions from acquiring any S or X locks on the table. Conversely, the first transaction to acquire an *intention shared* (IS) lock on a table prevents other transactions from acquiring any X locks on the table. The two-phase process allows the lock requests to be resolved in order, without blocking locks and corresponding operations that are compatible. For more information about this locking mechanism, see Section 15.5.1, “InnoDB Locking”.
See Also lock, lock mode, locking, transaction.

intention shared lock
See intention lock.

inverted index
A data structure optimized for document retrieval systems, used in the implementation of `InnoDB` **full-text search**. The `InnoDB` **FULLTEXT index**, implemented as an inverted index, records the position of each word within a document, rather than the location of a table row. A single column value (a document stored as a text string) is represented by many entries in the inverted index.
See Also full-text search, FULLTEXT index, ilist.

IOPS
Acronym for **I/O operations per second**. A common measurement for busy systems, particularly **OLTP** applications. If this value is near the maximum that the storage devices can handle, the application can become **disk-bound**, limiting **scalability**.
See Also disk-bound, OLTP, scalability.

isolation level
One of the foundations of database processing. Isolation is the **I** in the acronym **ACID**; the isolation level is the setting that fine-tunes the balance between performance and reliability, consistency, and reproducibility of results when multiple **transactions** are making changes and performing queries at the same time.

From highest amount of consistency and protection to the least, the isolation levels supported by InnoDB are: **SERIALIZABLE**, **REPEATABLE READ**, **READ COMMITTED**, and **READ UNCOMMITTED**.

With `InnoDB` tables, many users can keep the default isolation level (*REPEATABLE READ*) for all operations. Expert users might choose the **READ COMMITTED** level as they push the boundaries of scalability with **OLTP** processing, or during data warehousing operations where minor inconsistencies do not affect the aggregate results of large amounts of data. The levels on the edges (**SERIALIZABLE** and **READ UNCOMMITTED**) change the processing behavior to such an extent that they are rarely used.
See Also ACID, OLTP, READ COMMITTED, READ UNCOMMITTED, REPEATABLE READ, SERIALIZABLE, transaction.

J

join
A **query** that retrieves data from more than one table, by referencing columns in the tables that hold identical values. Ideally, these columns are part of an `InnoDB` **foreign key** relationship, which ensures **referential**

integrity and that the join columns are **indexed**. Often used to save space and improve query performance by replacing repeated strings with numeric IDs, in a **normalized** data design.
See Also foreign key, index, normalized, query, referential integrity.

K

KEY_BLOCK_SIZE
An option to specify the size of data pages within an `InnoDB` table that uses **compressed row format**. The default is 8 kilobytes. Lower values risk hitting internal limits that depend on the combination of row size and compression percentage.

For `MyISAM` tables, `KEY_BLOCK_SIZE` optionally specifies the size in bytes to use for index key blocks. The value is treated as a hint; a different size could be used if necessary. A `KEY_BLOCK_SIZE` value specified for an individual index definition overrides a table-level `KEY_BLOCK_SIZE` value.
See Also compressed row format.

L

latch
A lightweight structure used by `InnoDB` to implement a **lock** for its own internal memory structures, typically held for a brief time measured in milliseconds or microseconds. A general term that includes both **mutexes** (for exclusive access) and **rw-locks** (for shared access). Certain latches are the focus of `InnoDB` performance tuning. Statistics about latch use and contention are available through the **Performance Schema** interface.
See Also lock, locking, mutex, Performance Schema, rw-lock.

list
The `InnoDB` **buffer pool** is represented as a list of memory **pages**. The list is reordered as new pages are accessed and enter the buffer pool, as pages within the buffer pool are accessed again and are considered newer, and as pages that are not accessed for a long time are **evicted** from the buffer pool. The buffer pool is divided into **sublists**, and the replacement policy is a variation of the familiar **LRU** technique.
See Also buffer pool, eviction, LRU, page, sublist.

lock
The high-level notion of an object that controls access to a resource, such as a table, row, or internal data structure, as part of a **locking** strategy. For intensive performance tuning, you might delve into the actual structures that implement locks, such as **mutexes** and **latches**.
See Also latch, lock mode, locking, mutex.

lock escalation
An operation used in some database systems that converts many **row locks** into a single **table lock**, saving memory space but reducing concurrent access to the table. `InnoDB` uses a space-efficient representation for row locks, so that **lock** escalation is not needed.
See Also locking, row lock, table lock.

lock mode
A shared (S) **lock** allows a **transaction** to read a row. Multiple transactions can acquire an S lock on that same row at the same time.

An exclusive (X) lock allows a transaction to update or delete a row. No other transaction can acquire any kind of lock on that same row at the same time.

Intention locks apply to the table, and are used to indicate what kind of lock the transaction intends to acquire on rows in the table. Different transactions can acquire different kinds of intention locks on the same table, but the first transaction to acquire an intention exclusive (IX) lock on a table prevents other transactions from acquiring any S or X locks on the table. Conversely, the first transaction to acquire an intention shared (IS) lock on a table prevents other transactions from acquiring any X locks on the table. The two-phase process allows the lock requests to be resolved in order, without blocking locks and corresponding operations that are compatible.

See Also intention lock, lock, locking, transaction.

locking
The system of protecting a **transaction** from seeing or changing data that is being queried or changed by other transactions. The **locking** strategy must balance reliability and consistency of database operations (the principles of the **ACID** philosophy) against the performance needed for good **concurrency**. Fine-tuning the locking strategy often involves choosing an **isolation level** and ensuring all your database operations are safe and reliable for that isolation level.
See Also ACID, concurrency, isolation level, locking, transaction.

locking read
A `SELECT` statement that also performs a **locking** operation on an `InnoDB` table. Either `SELECT ... FOR UPDATE` or `SELECT ... LOCK IN SHARE MODE`. It has the potential to produce a **deadlock**, depending on the **isolation level** of the transaction. The opposite of a **non-locking read**. Not allowed for global tables in a **read-only transaction**.

`SELECT ... FOR SHARE` replaces `SELECT ... LOCK IN SHARE MODE` in MySQL 8.0.1, but `LOCK IN SHARE MODE` remains available for backward compatibility.

See Section 15.5.2.4, “Locking Reads”.
See Also deadlock, isolation level, locking, non-locking read, read-only transaction.

log
In the `InnoDB` context, “log” or “log files” typically refers to the **redo log** represented by the **ib_logfile*N*** files. Another type of `InnoDB` log is the **undo log**, which is a storage area that holds copies of data modified by active transactions.

Other kinds of logs that are important in MySQL are the **error log** (for diagnosing startup and runtime problems), **binary log** (for working with replication and performing point-in-time restores), the **general query log** (for diagnosing application problems), and the **slow query log** (for diagnosing performance problems).
See Also binary log, error log, general query log, ib_logfile, redo log, slow query log, undo log.

log buffer
The memory area that holds data to be written to the **log files** that make up the **redo log**. It is controlled by the `innodb_log_buffer_size` configuration option.
See Also log file, redo log.

log file
One of the **ib_logfile*N*** files that make up the **redo log**. Data is written to these files from the **log buffer** memory area.
See Also ib_logfile, log buffer, redo log.

log group
The set of files that make up the **redo log**, typically named `ib_logfile0` and `ib_logfile1`. (For that reason, sometimes referred to collectively as **ib_logfile**.)
See Also ib_logfile, redo log.

logical
A type of operation that involves high-level, abstract aspects such as tables, queries, indexes, and other SQL concepts. Typically, logical aspects are important to make database administration and application development convenient and usable. Contrast with **physical**.
See Also logical backup, physical.

logical backup
A **backup** that reproduces table structure and data, without copying the actual data files. For example, the **`mysqldump`** command produces a logical backup, because its output contains statements such as `CREATE TABLE` and `INSERT` that can re-create the data. Contrast with **physical backup**. A logical backup offers flexibility (for example, you could edit table definitions or insert statements before restoring), but can take substantially longer to **restore** than a physical backup.

See Also backup, mysqldump, physical backup, restore.

loose_
A prefix added to `InnoDB` configuration options after server **startup**, so any new configuration options not recognized by the current level of MySQL do not cause a startup failure. MySQL processes configuration options that start with this prefix, but gives a warning rather than a failure if the part after the prefix is not a recognized option.
See Also startup.

low-water mark
A value representing a lower limit, typically a threshold value at which some corrective action begins or becomes more aggressive. Contrast with **high-water mark**.
See Also high-water mark.

LRU
An acronym for “least recently used”, a common method for managing storage areas. The items that have not been used recently are **evicted** when space is needed to cache newer items. `InnoDB` uses the LRU mechanism by default to manage the **pages** within the **buffer pool**, but makes exceptions in cases where a page might be read only a single time, such as during a **full table scan**. This variation of the LRU algorithm is called the **midpoint insertion strategy**. For more information, see Section 15.6.3.1, “The InnoDB Buffer Pool”.
See Also buffer pool, eviction, full table scan, midpoint insertion strategy, page.

LSN
Acronym for “log sequence number”. This arbitrary, ever-increasing value represents a point in time corresponding to operations recorded in the **redo log**. (This point in time is regardless of **transaction** boundaries; it can fall in the middle of one or more transactions.) It is used internally by `InnoDB` during **crash recovery** and for managing the **buffer pool**.

Prior to MySQL 5.6.3, the LSN was a 4-byte unsigned integer. The LSN became an 8-byte unsigned integer in MySQL 5.6.3 when the redo log file size limit increased from 4GB to 512GB, as additional bytes were required to store extra size information. Applications built on MySQL 5.6.3 or later that use LSN values should use 64-bit rather than 32-bit variables to store and compare LSN values.

In the **MySQL Enterprise Backup** product, you can specify an LSN to represent the point in time from which to take an **incremental backup**. The relevant LSN is displayed by the output of the `mysqlbackup` command. Once you have the LSN corresponding to the time of a full backup, you can specify that value to take a subsequent incremental backup, whose output contains another LSN for the next incremental backup.
See Also buffer pool, crash recovery, incremental backup, MySQL Enterprise Backup, redo log, transaction.

M

.MRG file
A file containing references to other tables, used by the `MERGE` storage engine. Files with this extension are always included in backups produced by the `mysqlbackup` command of the **MySQL Enterprise Backup** product.
See Also MySQL Enterprise Backup, mysqlbackup command.

.MYD file
A file that MySQL uses to store data for a `MyISAM` table.
See Also .MYI file, MySQL Enterprise Backup, mysqlbackup command.

.MYI file
A file that MySQL uses to store indexes for a `MyISAM` table.
See Also .MYD file, MySQL Enterprise Backup, mysqlbackup command.

master server
Frequently shortened to “master”. A database server machine in a **replication** scenario that processes the initial insert, update, and delete requests for data. These changes are propagated to, and repeated on, other servers known as **slave servers**.

See Also replication, slave server.

master thread
An InnoDB **thread** that performs various tasks in the background. Most of these tasks are I/O related, such as writing changes from the **change buffer** to the appropriate secondary indexes.

To improve **concurrency**, sometimes actions are moved from the master thread to separate background threads. For example, in MySQL 5.6 and higher, **dirty pages** are **flushed** from the **buffer pool** by the **page cleaner** thread rather than the master thread.
See Also buffer pool, change buffer, concurrency, dirty page, flush, page cleaner, thread.

MDL
Acronym for “metadata lock”.
See Also metadata lock.

memcached
A popular component of many MySQL and **NoSQL** software stacks, allowing fast reads and writes for single values and caching the results entirely in memory. Traditionally, applications required extra logic to write the same data to a MySQL database for permanent storage, or to read data from a MySQL database when it was not cached yet in memory. Now, applications can use the simple memcached protocol, supported by client libraries for many languages, to communicate directly with MySQL servers using InnoDB or NDB tables. These NoSQL interfaces to MySQL tables allow applications to achieve higher read and write performance than by issuing SQL statements directly, and can simplify application logic and deployment configurations for systems that already incorporate memcached for in-memory caching.

The memcached interface to InnoDB tables is available in MySQL 5.6 and higher; see Section 15.19, “InnoDB memcached Plugin” for details. The memcached interface to NDB tables is available in NDB Cluster 7.2 and later; see http://dev.mysql.com/doc/ndbapi/en/ndbmemcache.html for details.
See Also InnoDB, NoSQL.

merge
To apply changes to data cached in memory, such as when a page is brought into the **buffer pool**, and any applicable changes recorded in the **change buffer** are incorporated into the page in the buffer pool. The updated data is eventually written to the **tablespace** by the **flush** mechanism.
See Also buffer pool, change buffer, flush, tablespace.

metadata lock
A type of **lock** that prevents **DDL** operations on a table that is being used at the same time by another **transaction**. For details, see Section 8.11.4, “Metadata Locking”.

Enhancements to **online** operations, particularly in MySQL 5.6 and higher, are focused on reducing the amount of metadata locking. The objective is for DDL operations that do not change the table structure (such as CREATE INDEX and DROP INDEX for InnoDB tables) to proceed while the table is being queried, updated, and so on by other transactions.
See Also DDL, lock, online, transaction.

metrics counter
A feature implemented by the INNODB_METRICS table in the **INFORMATION_SCEMA**, in MySQL 5.6 and higher. You can query **counts** and totals for low-level InnoDB operations, and use the results for performance tuning in combination with data from the **Performance Schema**.
See Also counter, INFORMATION_SCHEMA, Performance Schema.

midpoint insertion strategy
The technique of initially bringing **pages** into the InnoDB **buffer pool** not at the “newest” end of the list, but instead somewhere in the middle. The exact location of this point can vary, based on the setting of the innodb_old_blocks_pct option. The intent is that pages that are only read once, such as during a **full table scan**, can be aged out of the buffer pool sooner than with a strict **LRU** algorithm. For more information, see Section 15.6.3.1, “The InnoDB Buffer Pool”.
See Also buffer pool, full table scan, LRU, page.

mini-transaction
An internal phase of InnoDB processing, when making changes at the **physical** level to internal data structures during **DML** operations. A mini-transaction (mtr) has no notion of **rollback**; multiple mini-transactions can occur within a single **transaction**. Mini-transactions write information to the **redo log** that is used during **crash recovery**. A mini-transaction can also happen outside the context of a regular transaction, for example during **purge** processing by background threads.
See Also commit, crash recovery, DML, physical, purge, redo log, rollback, transaction.

mixed-mode insert
An INSERT statement where **auto-increment** values are specified for some but not all of the new rows. For example, a multi-value INSERT could specify a value for the auto-increment column in some cases and NULL in other cases. InnoDB generates auto-increment values for the rows where the column value was specified as NULL. Another example is an INSERT ... ON DUPLICATE KEY UPDATE statement, where auto-increment values might be generated but not used, for any duplicate rows that are processed as UPDATE rather than INSERT statements.

Can cause consistency issues between **master** and **slave** servers in a **replication** configuration. Can require adjusting the value of the **innodb_autoinc_lock_mode** configuration option.
See Also auto-increment, innodb_autoinc_lock_mode, master server, replication, slave server.

mtr
See mini-transaction.

multi-core
A type of processor that can take advantage of multi-threaded programs, such as the MySQL server.

multiversion concurrency control
See MVCC.

mutex
Informal abbreviation for “mutex variable”. (Mutex itself is short for “mutual exclusion”.) The low-level object that InnoDB uses to represent and enforce exclusive-access **locks** to internal in-memory data structures. Once the lock is acquired, any other process, thread, and so on is prevented from acquiring the same lock. Contrast with **rw-locks**, which InnoDB uses to represent and enforce shared-access **locks** to internal in-memory data structures. Mutexes and rw-locks are known collectively as **latches**.
See Also latch, lock, Performance Schema, Pthreads, rw-lock.

MVCC
Acronym for “multiversion concurrency control”. This technique lets InnoDB **transactions** with certain **isolation levels** perform **consistent read** operations; that is, to query rows that are being updated by other transactions, and see the values from before those updates occurred. This is a powerful technique to increase **concurrency**, by allowing queries to proceed without waiting due to **locks** held by the other transactions.

This technique is not universal in the database world. Some other database products, and some other MySQL storage engines, do not support it.
See Also ACID, concurrency, consistent read, isolation level, lock, transaction.

my.cnf
The name, on Unix or Linux systems, of the MySQL **option file**.
See Also my.ini, option file.

my.ini
The name, on Windows systems, of the MySQL **option file**.
See Also my.cnf, option file.

mysql
The mysql program is the command-line interpreter for the MySQL database. It processes **SQL** statements, and also MySQL-specific commands such as SHOW TABLES, by passing requests to the mysqld daemon.
See Also mysqld, SQL.

MySQL Enterprise Backup
: A licensed product that performs **hot backups** of MySQL databases. It offers the most efficiency and flexibility when backing up `InnoDB` tables, but can also back up `MyISAM` and other kinds of tables.
See Also hot backup, InnoDB.

mysqlbackup command
: A command-line tool of the **MySQL Enterprise Backup** product. It performs a **hot backup** operation for `InnoDB` tables, and a warm backup for `MyISAM` and other kinds of tables. See Section 29.2, "MySQL Enterprise Backup Overview" for more information about this command.
See Also hot backup, MySQL Enterprise Backup, warm backup.

mysqld
: The `mysqld` program is the database engine for the MySQL database. It runs as a Unix daemon or Windows service, constantly waiting for requests and performing maintenance work in the background.
See Also mysql.

mysqldump
: A command that performs a **logical backup** of some combination of databases, tables, and table data. The results are SQL statements that reproduce the original schema objects, data, or both. For substantial amounts of data, a **physical backup** solution such as **MySQL Enterprise Backup** is faster, particularly for the **restore** operation.
See Also logical backup, MySQL Enterprise Backup, physical backup, restore.

N

natural key
: An indexed column, typically a **primary key**, where the values have some real-world significance. Usually advised against because:

  - If the value should ever change, there is potentially a lot of index maintenance to re-sort the **clustered index** and update the copies of the primary key value that are repeated in each **secondary index**.
  - Even seemingly stable values can change in unpredictable ways that are difficult to represent correctly in the database. For example, one country can change into two or several, making the original country code obsolete. Or, rules about unique values might have exceptions. For example, even if taxpayer IDs are intended to be unique to a single person, a database might have to handle records that violate that rule, such as in cases of identity theft. Taxpayer IDs and other sensitive ID numbers also make poor primary keys, because they may need to be secured, encrypted, and otherwise treated differently than other columns.

  Thus, it is typically better to use arbitrary numeric values to form a **synthetic key**, for example using an **auto-increment** column.
  See Also auto-increment, clustered index, primary key, secondary index, synthetic key.

neighbor page
: Any **page** in the same **extent** as a particular page. When a page is selected to be **flushed**, any neighbor pages that are **dirty** are typically flushed as well, as an I/O optimization for traditional hard disks. In MySQL 5.6 and up, this behavior can be controlled by the configuration variable `innodb_flush_neighbors`; you might turn that setting off for SSD drives, which do not have the same overhead for writing smaller batches of data at random locations.
See Also dirty page, extent, flush, page.

next-key lock
: A combination of a **record lock** on the index record and a gap lock on the gap before the index record.
See Also gap lock, locking, record lock.

non-blocking I/O
: An industry term that means the same as **asynchronous I/O**.

See Also asynchronous I/O.

non-locking read
A **query** that does not use the `SELECT ... FOR UPDATE` or `SELECT ... LOCK IN SHARE MODE` clauses. The only kind of query allowed for global tables in a **read-only transaction**. The opposite of a **locking read**. See Section 15.5.2.3, "Consistent Nonlocking Reads".

`SELECT ... FOR SHARE` replaces `SELECT ... LOCK IN SHARE MODE` in MySQL 8.0.1, but `LOCK IN SHARE MODE` remains available for backward compatibility.
See Also locking read, query, read-only transaction.

non-repeatable read
The situation when a query retrieves data, and a later query within the same **transaction** retrieves what should be the same data, but the queries return different results (changed by another transaction committing in the meantime).

This kind of operation goes against the **ACID** principle of database design. Within a transaction, data should be consistent, with predictable and stable relationships.

Among different **isolation levels**, non-repeatable reads are prevented by the **serializable read** and **repeatable read** levels, and allowed by the **consistent read**, and **read uncommitted** levels.
See Also ACID, consistent read, isolation level, READ UNCOMMITTED, REPEATABLE READ, SERIALIZABLE, transaction.

normalized
A database design strategy where data is split into multiple tables, and duplicate values condensed into single rows represented by an ID, to avoid storing, querying, and updating redundant or lengthy values. It is typically used in **OLTP** applications.

For example, an address might be given a unique ID, so that a census database could represent the relationship **lives at this address** by associating that ID with each member of a family, rather than storing multiple copies of a complex value such as **123 Main Street, Anytown, USA**.

For another example, although a simple address book application might store each phone number in the same table as a person's name and address, a phone company database might give each phone number a special ID, and store the numbers and IDs in a separate table. This normalized representation could simplify large-scale updates when area codes split apart.

Normalization is not always recommended. Data that is primarily queried, and only updated by deleting entirely and reloading, is often kept in fewer, larger tables with redundant copies of duplicate values. This data representation is referred to as **denormalized**, and is frequently found in data warehousing applications.
See Also denormalized, foreign key, OLTP, relational.

NoSQL
A broad term for a set of data access technologies that do not use the **SQL** language as their primary mechanism for reading and writing data. Some NoSQL technologies act as key-value stores, only accepting single-value reads and writes; some relax the restrictions of the **ACID** methodology; still others do not require a pre-planned **schema**. MySQL users can combine NoSQL-style processing for speed and simplicity with SQL operations for flexibility and convenience, by using the **memcached** API to directly access some kinds of MySQL tables. The `memcached` interface to `InnoDB` tables is available in MySQL 5.6 and higher; see Section 15.19, "InnoDB memcached Plugin" for details. The `memcached` interface to `NDB` tables is available in NDB Cluster 7.2 and later; see ndbmemcache—Memcache API for NDB Cluster.
See Also ACID, InnoDB, memcached, schema, SQL.

NOT NULL constraint
A type of **constraint** that specifies that a **column** cannot contain any **NULL** values. It helps to preserve **referential integrity**, as the database server can identify data with erroneous missing values. It also helps in the arithmetic involved in query optimization, allowing the optimizer to predict the number of entries in an index on that column.
See Also column, constraint, NULL, primary key, referential integrity.

NULL
A special value in **SQL**, indicating the absence of data. Any arithmetic operation or equality test involving a `NULL` value, in turn produces a `NULL` result. (Thus it is similar to the IEEE floating-point concept of NaN, “not a number”.) Any aggregate calculation such as `AVG()` ignores rows with `NULL` values, when determining how many rows to divide by. The only test that works with `NULL` values uses the SQL idioms `IS NULL` or `IS NOT NULL`.

`NULL` values play a part in **index** operations, because for performance a database must minimize the overhead of keeping track of missing data values. Typically, `NULL` values are not stored in an index, because a query that tests an indexed column using a standard comparison operator could never match a row with a `NULL` value for that column. For the same reason, unique indexes do not prevent `NULL` values; those values simply are not represented in the index. Declaring a `NOT NULL` constraint on a column provides reassurance that there are no rows left out of the index, allowing for better query optimization (accurate counting of rows and estimation of whether to use the index).

Because the **primary key** must be able to uniquely identify every row in the table, a single-column primary key cannot contain any `NULL` values, and a multi-column primary key cannot contain any rows with `NULL` values in all columns.

Although the Oracle database allows a `NULL` value to be concatenated with a string, `InnoDB` treats the result of such an operation as `NULL`.
See Also index, primary key, SQL.

O

.OPT file
A file containing database configuration information. Files with this extension are included in backups produced by the `mysqlbackup` command of the **MySQL Enterprise Backup** product.
See Also MySQL Enterprise Backup, mysqlbackup command.

off-page column
A column containing variable-length data (such as `BLOB` and `VARCHAR`) that is too long to fit on a **B-tree** page. The data is stored in **overflow pages**. The **DYNAMIC** row format is more efficient for such storage than the older **COMPACT** row format.
See Also B-tree, compact row format, dynamic row format, overflow page.

OLTP
Acronym for “Online Transaction Processing”. A database system, or a database application, that runs a workload with many **transactions**, with frequent writes as well as reads, typically affecting small amounts of data at a time. For example, an airline reservation system or an application that processes bank deposits. The data might be organized in **normalized** form for a balance between **DML** (insert/update/delete) efficiency and **query** efficiency. Contrast with **data warehouse**.

With its **row-level locking** and **transactional** capability, **InnoDB** is the ideal storage engine for MySQL tables used in OLTP applications.
See Also data warehouse, DML, InnoDB, query, row lock, transaction.

online
A type of operation that involves no downtime, blocking, or restricted operation for the database. Typically applied to **DDL**. Operations that shorten the periods of restricted operation, such as **fast index creation**, have evolved into a wider set of **online DDL** operations in MySQL 5.6.

In the context of backups, a **hot backup** is an online operation and a **warm backup** is partially an online operation.
See Also DDL, Fast Index Creation, hot backup, online DDL, warm backup.

online DDL
A feature that improves the performance, concurrency, and availability of `InnoDB` tables during **DDL** (primarily `ALTER TABLE`) operations. See Section 15.12, “InnoDB and Online DDL” for details.

The details vary according to the type of operation. In some cases, the table can be modified concurrently while the `ALTER TABLE` is in progress. The operation might be able to be performed without a table copy, or using a specially optimized type of table copy. Space usage is controlled by the `innodb_online_alter_log_max_size` configuration option.

This feature is an enhancement of the **Fast Index Creation** feature in MySQL 5.5.
See Also DDL, Fast Index Creation, online.

optimistic
A methodology that guides low-level implementation decisions for a relational database system. The requirements of performance and **concurrency** in a relational database mean that operations must be started or dispatched quickly. The requirements of consistency and **referential integrity** mean that any operation could fail: a transaction might be rolled back, a **DML** operation could violate a constraint, a request for a lock could cause a deadlock, a network error could cause a timeout. An optimistic strategy is one that assumes most requests or attempts will succeed, so that relatively little work is done to prepare for the failure case. When this assumption is true, the database does little unnecessary work; when requests do fail, extra work must be done to clean up and undo changes.

`InnoDB` uses optimistic strategies for operations such as **locking** and **commits**. For example, data changed by a transaction can be written to the data files before the commit occurs, making the commit itself very fast, but requiring more work to undo the changes if the transaction is rolled back.

The opposite of an optimistic strategy is a **pessimistic** one, where a system is optimized to deal with operations that are unreliable and frequently unsuccessful. This methodology is rare in a database system, because so much care goes into choosing reliable hardware, networks, and algorithms.
See Also commit, concurrency, DML, locking, pessimistic, referential integrity.

optimizer
The MySQL component that determines the best **indexes** and **join** order to use for a **query**, based on characteristics and data distribution of the relevant **tables**.
See Also index, join, query, table.

option
A configuration parameter for MySQL, either stored in the **option file** or passed on the command line.

For the **options** that apply to **InnoDB** tables, each option name starts with the prefix `innodb_`.
See Also InnoDB, option, option file.

option file
The file that holds the configuration **options** for the MySQL instance. Traditionally, on Linux and Unix this file is named `my.cnf`, and on Windows it is named `my.ini`.
See Also configuration file, my.cnf, my.ini, option.

overflow page
Separately allocated disk **pages** that hold variable-length columns (such as `BLOB` and `VARCHAR`) that are too long to fit on a **B-tree** page. The associated columns are known as **off-page columns**.
See Also B-tree, off-page column, page.

P

.PAR file
A file containing partition definitions. Files with this extension are always included in backups produced by the `mysqlbackup` command of the **MySQL Enterprise Backup** product.

With the introduction of native partitioning support for `InnoDB` tables in MySQL 5.7.6, `.PAR` files are no longer created for partitioned `InnoDB` tables.
See Also MySQL Enterprise Backup, mysqlbackup command.

page
A unit representing how much data `InnoDB` transfers at any one time between disk (the **data files**) and memory (the **buffer pool**). A page can contain one or more **rows**, depending on how much data is in each

row. If a row does not fit entirely into a single page, `InnoDB` sets up additional pointer-style data structures so that the information about the row can be stored in one page.

One way to fit more data in each page is to use **compressed row format**. For tables that use BLOBs or large text fields, **compact row format** allows those large columns to be stored separately from the rest of the row, reducing I/O overhead and memory usage for queries that do not reference those columns.

When `InnoDB` reads or writes sets of pages as a batch to increase I/O throughput, it reads or writes an **extent** at a time.

All the `InnoDB` disk data structures within a MySQL instance share the same **page size**.
See Also buffer pool, compact row format, compressed row format, data files, extent, page size, row.

page cleaner
An `InnoDB` background **thread** that **flushes dirty pages** from the **buffer pool**. Prior to MySQL 5.6, this activity was performed by the **master thread**. The number of page cleaner threads is controlled by the `innodb_page_cleaners` configuration option, introduced in MySQL 5.7.4.
See Also buffer pool, dirty page, flush, master thread, thread.

page size
For releases up to and including MySQL 5.5, the size of each `InnoDB` **page** is fixed at 16 kilobytes. This value represents a balance: large enough to hold the data for most rows, yet small enough to minimize the performance overhead of transferring unneeded data to memory. Other values are not tested or supported.

Starting in MySQL 5.6, the page size for an `InnoDB` **instance** can be either 4KB, 8KB, or 16KB, controlled by the `innodb_page_size` configuration option. As of MySQL 5.7.6, `InnoDB` also supports 32KB and 64KB page sizes. For 32KB and 64KB page sizes, `ROW_FORMAT=COMPRESSED` is not supported and the maximum record size is 16KB.

Page size is set when creating the MySQL instance, and it remains constant afterward. The same page size applies to all `InnoDB` **tablespaces**, including the **system tablespace**, **file-per-table** tablespaces, and **general tablespaces**.

Smaller page sizes can help performance with storage devices that use small block sizes, particularly for **SSD** devices in **disk-bound** workloads, such as for **OLTP** applications. As individual rows are updated, less data is copied into memory, written to disk, reorganized, locked, and so on.
See Also disk-bound, file-per-table, general tablespace, instance, OLTP, page, SSD, system tablespace, tablespace.

parent table
The table in a **foreign key** relationship that holds the initial column values pointed to from the **child table**. The consequences of deleting, or updating rows in the parent table depend on the `ON UPDATE` and `ON DELETE` clauses in the foreign key definition. Rows with corresponding values in the child table could be automatically deleted or updated in turn, or those columns could be set to `NULL`, or the operation could be prevented.
See Also child table, foreign key.

partial backup
A **backup** that contains some of the **tables** in a MySQL database, or some of the databases in a MySQL instance. Contrast with **full backup**.
See Also backup, full backup, table.

partial index
An **index** that represents only part of a column value, typically the first N characters (the **prefix**) of a long `VARCHAR` value.
See Also index, index prefix.

Performance Schema
The `performance_schema` schema, in MySQL 5.5 and up, presents a set of tables that you can query to get detailed information about the performance characteristics of many internal parts of the MySQL server.
See Chapter 25, *MySQL Performance Schema*.

See Also INFORMATION_SCHEMA, latch, mutex, rw-lock.

persistent statistics
A feature in MySQL 5.6 that stores **index** statistics for `InnoDB` **tables** on disk, providing better **plan stability** for **queries**. For more information, see Section 15.6.11.1, “Configuring Persistent Optimizer Statistics Parameters”.
See Also index, optimizer, plan stability, query, table.

pessimistic
A methodology that sacrifices performance or concurrency in favor of safety. It is appropriate if a high proportion of requests or attempts might fail, or if the consequences of a failed request are severe. `InnoDB` uses what is known as a pessimistic **locking** strategy, to minimize the chance of **deadlocks**. At the application level, you might avoid deadlocks by using a pessimistic strategy of acquiring all locks needed by a transaction at the very beginning.

Many built-in database mechanisms use the opposite **optimistic** methodology.
See Also deadlock, locking, optimistic.

phantom
A row that appears in the result set of a query, but not in the result set of an earlier query. For example, if a query is run twice within a **transaction**, and in the meantime, another transaction commits after inserting a new row or updating a row so that it matches the `WHERE` clause of the query.

This occurrence is known as a phantom read. It is harder to guard against than a **non-repeatable read**, because locking all the rows from the first query result set does not prevent the changes that cause the phantom to appear.

Among different **isolation levels**, phantom reads are prevented by the **serializable read** level, and allowed by the **repeatable read**, **consistent read**, and **read uncommitted** levels.
See Also consistent read, isolation level, non-repeatable read, READ UNCOMMITTED, REPEATABLE READ, SERIALIZABLE, transaction.

physical
A type of operation that involves hardware-related aspects such as disk blocks, memory pages, files, bits, disk reads, and so on. Typically, physical aspects are important during expert-level performance tuning and problem diagnosis. Contrast with **logical**.
See Also logical, physical backup.

physical backup
A **backup** that copies the actual data files. For example, the **`mysqlbackup`** command of the **MySQL Enterprise Backup** product produces a physical backup, because its output contains data files that can be used directly by the `mysqld` server, resulting in a faster **restore** operation. Contrast with **logical backup**.
See Also backup, logical backup, MySQL Enterprise Backup, restore.

PITR
Acronym for **point-in-time recovery**.
See Also point-in-time recovery.

plan stability
A property of a **query execution plan**, where the optimizer makes the same choices each time for a given **query**, so that performance is consistent and predictable.
See Also query, query execution plan.

point-in-time recovery
The process of restoring a **backup** to recreate the state of the database at a specific date and time. Commonly abbreviated “PITR”. Because it is unlikely that the specified time corresponds exactly to the time of a backup, this technique usually requires a combination of a **physical backup** and a **logical backup**. For example, with the **MySQL Enterprise Backup** product, you restore the last backup that you took before the specified point in time, then replay changes from the **binary log** between the time of the backup and the PITR time.
See Also backup, binary log, logical backup, MySQL Enterprise Backup, physical backup.

prefix
See index prefix.

prepared backup
A set of backup files, produced by the **MySQL Enterprise Backup** product, after all the stages of applying **binary logs** and **incremental backups** are finished. The resulting files are ready to be **restored**. Prior to the apply steps, the files are known as a **raw backup**.
See Also binary log, hot backup, incremental backup, MySQL Enterprise Backup, raw backup, restore.

primary key
A set of columns—and by implication, the index based on this set of columns—that can uniquely identify every row in a table. As such, it must be a unique index that does not contain any `NULL` values.

`InnoDB` requires that every table has such an index (also called the **clustered index** or **cluster index**), and organizes the table storage based on the column values of the primary key.

When choosing primary key values, consider using arbitrary values (a **synthetic key**) rather than relying on values derived from some other source (a **natural key**).
See Also clustered index, index, natural key, synthetic key.

process
An instance of an executing program. The operating system switches between multiple running processes, allowing for a certain degree of **concurrency**. On most operating systems, processes can contain multiple **threads** of execution that share resources. Context-switching between threads is faster than the equivalent switching between processes.
See Also concurrency, thread.

pseudo-record
An artificial record in an index, used for **locking** key values or ranges that do not currently exist.
See Also infimum record, locking, supremum record.

Pthreads
The POSIX threads standard, which defines an API for threading and locking operations on Unix and Linux systems. On Unix and Linux systems, `InnoDB` uses this implementation for **mutexes**.
See Also mutex.

purge
A type of garbage collection performed by one or more separate background threads (controlled by `innodb_purge_threads`) that runs on a periodic schedule. Purge parses and processes **undo log** pages from the **history list** for the purpose of removing clustered and secondary index records that were marked for deletion (by previous `DELETE` statements) and are no longer required for **MVCC** or **rollback**. Purge frees undo log pages from the history list after processing them.
See Also history list, MVCC, rollback, undo log.

purge buffering
The technique of storing changes to secondary index pages, resulting from `DELETE` operations, in the **change buffer** rather than writing the changes immediately, so that the physical writes can be performed to minimize random I/O. (Because delete operations are a two-step process, this operation buffers the write that normally purges an index record that was previously marked for deletion.) It is one of the types of **change buffering**; the others are **insert buffering** and **delete buffering**.
See Also change buffer, change buffering, delete buffering, insert buffer, insert buffering.

purge lag
Another name for the `InnoDB` **history list**. Related to the `innodb_max_purge_lag` configuration option.
See Also history list, purge.

purge thread
A **thread** within the `InnoDB` process that is dedicated to performing the periodic **purge** operation. In MySQL 5.6 and higher, multiple purge threads are enabled by the `innodb_purge_threads` configuration option.
See Also purge, thread.

Q

query
: In **SQL**, an operation that reads information from one or more **tables**. Depending on the organization of data and the parameters of the query, the lookup might be optimized by consulting an **index**. If multiple tables are involved, the query is known as a **join**.

  For historical reasons, sometimes discussions of internal processing for statements use “query” in a broader sense, including other types of MySQL statements such as **DDL** and **DML** statements.
  See Also DDL, DML, index, join, SQL, table.

query execution plan
: The set of decisions made by the optimizer about how to perform a **query** most efficiently, including which **index** or indexes to use, and the order in which to **join** tables. **Plan stability** involves the same choices being made consistently for a given query.
  See Also index, join, plan stability, query.

query log
: See general query log.

quiesce
: To reduce the amount of database activity, often in preparation for an operation such as an `ALTER TABLE`, a **backup**, or a **shutdown**. Might or might not involve doing as much **flushing** as possible, so that **InnoDB** does not continue doing background I/O.

  In MySQL 5.6 and higher, the syntax `FLUSH TABLES ... FOR EXPORT` writes some data to disk for `InnoDB` tables that make it simpler to back up those tables by copying the data files.
  See Also backup, flush, InnoDB, shutdown.

R

R-tree
: A tree data structure used for spatial indexing of multi-dimensional data such as geographical coordinates, rectangles or polygons.
  See Also B-tree.

RAID
: Acronym for “Redundant Array of Inexpensive Drives”. Spreading I/O operations across multiple drives enables greater **concurrency** at the hardware level, and improves the efficiency of low-level write operations that otherwise would be performed in sequence.
  See Also concurrency.

random dive
: A technique for quickly estimating the number of different values in a column (the column's **cardinality**). `InnoDB` samples pages at random from the index and uses that data to estimate the number of different values.
  See Also cardinality.

raw backup
: The initial set of backup files produced by the **MySQL Enterprise Backup** product, before the changes reflected in the **binary log** and any **incremental backups** are applied. At this stage, the files are not ready to **restore**. After these changes are applied, the files are known as a **prepared backup**.
  See Also binary log, hot backup, ibbackup_logfile, incremental backup, MySQL Enterprise Backup, prepared backup, restore.

READ COMMITTED
: An **isolation level** that uses a **locking** strategy that relaxes some of the protection between **transactions**, in the interest of performance. Transactions cannot see uncommitted data from other transactions, but they can see data that is committed by another transaction after the current transaction started. Thus, a transaction

never sees any bad data, but the data that it does see may depend to some extent on the timing of other transactions.

When a transaction with this isolation level performs `UPDATE ... WHERE` or `DELETE ... WHERE` operations, other transactions might have to wait. The transaction can perform `SELECT ... FOR UPDATE`, and `LOCK IN SHARE MODE` operations without making other transactions wait.

`SELECT ... FOR SHARE` replaces `SELECT ... LOCK IN SHARE MODE` in MySQL 8.0.1, but `LOCK IN SHARE MODE` remains available for backward compatibility.
See Also ACID, isolation level, locking, REPEATABLE READ, SERIALIZABLE, transaction.

read phenomena
Phenomena such as **dirty reads**, **non-repeatable reads**, and **phantom** reads which can occur when a transaction reads data that another transaction has modified.
See Also dirty read, non-repeatable read, phantom.

READ UNCOMMITTED
The **isolation level** that provides the least amount of protection between transactions. Queries employ a **locking** strategy that allows them to proceed in situations where they would normally wait for another transaction. However, this extra performance comes at the cost of less reliable results, including data that has been changed by other transactions and not committed yet (known as **dirty read**). Use this isolation level with great caution, and be aware that the results might not be consistent or reproducible, depending on what other transactions are doing at the same time. Typically, transactions with this isolation level only do queries, not insert, update, or delete operations.
See Also ACID, dirty read, isolation level, locking, transaction.

read view
An internal snapshot used by the **MVCC** mechanism of `InnoDB`. Certain **transactions**, depending on their **isolation level**, see the data values as they were at the time the transaction (or in some cases, the statement) started. Isolation levels that use a read view are **REPEATABLE READ**, **READ COMMITTED**, and **READ UNCOMMITTED**.
See Also isolation level, MVCC, READ COMMITTED, READ UNCOMMITTED, REPEATABLE READ, transaction.

read-ahead
A type of I/O request that prefetches a group of **pages** (an entire **extent**) into the **buffer pool** asynchronously, in anticipation that these pages will be needed soon. The linear read-ahead technique prefetches all the pages of one extent based on access patterns for pages in the preceding extent. The random read-ahead technique prefetches all the pages for an extent once a certain number of pages from the same extent are in the buffer pool. Random read-ahead is not part of MySQL 5.5, but is re-introduced in MySQL 5.6 under the control of the `innodb_random_read_ahead` configuration option.
See Also buffer pool, extent, page.

read-only transaction
A type of **transaction** that can be optimized for `InnoDB` tables by eliminating some of the bookkeeping involved with creating a **read view** for each transaction. Can only perform **non-locking read** queries. It can be started explicitly with the syntax `START TRANSACTION READ ONLY`, or automatically under certain conditions. See Section 8.5.3, “Optimizing InnoDB Read-Only Transactions” for details.
See Also non-locking read, read view, transaction.

record lock
A lock on an index record. For example, `SELECT c1 FROM t WHERE c1 = 10 FOR UPDATE;` prevents any other transaction from inserting, updating, or deleting rows where the value of `t.c1` is 10. Contrast with **gap lock** and **next-key lock**.
See Also gap lock, lock, next-key lock.

redo
The data, in units of records, recorded in the **redo log** when DML statements make changes to `InnoDB` tables. It is used during **crash recovery** to correct data written by incomplete **transactions**. The ever-increasing **LSN** value represents the cumulative amount of redo data that has passed through the redo log.

See Also crash recovery, DML, LSN, redo log, transaction.

redo log

A disk-based data structure used during **crash recovery**, to correct data written by incomplete **transactions**. During normal operation, it encodes requests to change `InnoDB` table data, which result from SQL statements or low-level API calls through NoSQL interfaces. Modifications that did not finish updating the **data files** before an unexpected **shutdown** are replayed automatically.

The redo log is physically represented as a set of files, typically named `ib_logfile0` and `ib_logfile1`. The data in the redo log is encoded in terms of records affected; this data is collectively referred to as **redo**. The passage of data through the redo logs is represented by the ever-increasing **LSN** value. The original 4GB limit on maximum size for the redo log is raised to 512GB in MySQL 5.6.3.

The disk layout of the redo log is influenced by the configuration options `innodb_log_file_size`, `innodb_log_group_home_dir`, and (rarely) `innodb_log_files_in_group`. The performance of redo log operations is also affected by the **log buffer**, which is controlled by the `innodb_log_buffer_size` configuration option.
See Also crash recovery, data files, ib_logfile, log buffer, LSN, redo, shutdown, transaction.

redundant row format

The oldest `InnoDB` **row format**. Prior to MySQL 5.0.3, it was the only row format available in `InnoDB`. From MySQL 5.0.3 to MySQL 5.7.8, the default row format is **COMPACT**. As of MySQL 5.7.9, the default row format is defined by the `innodb_default_row_format` configuration option, which has a default setting of **DYNAMIC**. You can still specify the **REDUNDANT** row format for compatibility with older `InnoDB` tables.

For more information, see Section 15.10.4, “COMPACT and REDUNDANT Row Formats”.
See Also compact row format, dynamic row format, row format.

referential integrity

The technique of maintaining data always in a consistent format, part of the **ACID** philosophy. In particular, data in different tables is kept consistent through the use of **foreign key constraints**, which can prevent changes from happening or automatically propagate those changes to all related tables. Related mechanisms include the **unique constraint**, which prevents duplicate values from being inserted by mistake, and the **NOT NULL constraint**, which prevents blank values from being inserted by mistake.
See Also ACID, FOREIGN KEY constraint, NOT NULL constraint, unique constraint.

relational

An important aspect of modern database systems. The database server encodes and enforces relationships such as one-to-one, one-to-many, many-to-one, and uniqueness. For example, a person might have zero, one, or many phone numbers in an address database; a single phone number might be associated with several family members. In a financial database, a person might be required to have exactly one taxpayer ID, and any taxpayer ID could only be associated with one person.

The database server can use these relationships to prevent bad data from being inserted, and to find efficient ways to look up information. For example, if a value is declared to be unique, the server can stop searching as soon as the first match is found, and it can reject attempts to insert a second copy of the same value.

At the database level, these relationships are expressed through SQL features such as **columns** within a table, unique and `NOT NULL` **constraints**, **foreign keys**, and different kinds of join operations. Complex relationships typically involve data split between more than one table. Often, the data is **normalized**, so that duplicate values in one-to-many relationships are stored only once.

In a mathematical context, the relations within a database are derived from set theory. For example, the `OR` and `AND` operators of a `WHERE` clause represent the notions of union and intersection.
See Also ACID, column, constraint, foreign key, normalized.

relevance

In the **full-text search** feature, a number signifying the similarity between the search string and the data in the **FULLTEXT index**. For example, when you search for a single word, that word is typically more relevant for a row where it occurs several times in the text than a row where it appears only once.
See Also full-text search, FULLTEXT index.

REPEATABLE READ
: The default **isolation level** for `InnoDB`. It prevents any rows that are queried from being changed by other **transactions**, thus blocking **non-repeatable reads** but not **phantom** reads. It uses a moderately strict **locking** strategy so that all queries within a transaction see data from the same snapshot, that is, the data as it was at the time the transaction started.

  When a transaction with this isolation level performs `UPDATE ... WHERE`, `DELETE ... WHERE`, `SELECT ... FOR UPDATE`, and `LOCK IN SHARE MODE` operations, other transactions might have to wait.

  `SELECT ... FOR SHARE` replaces `SELECT ... LOCK IN SHARE MODE` in MySQL 8.0.1, but `LOCK IN SHARE MODE` remains available for backward compatibility.
  See Also ACID, consistent read, isolation level, locking, phantom, transaction.

replication
: The practice of sending changes from a **master database**, to one or more **slave databases**, so that all databases have the same data. This technique has a wide range of uses, such as load-balancing for better scalability, disaster recovery, and testing software upgrades and configuration changes. The changes can be sent between the database by methods called **row-based replication** and **statement-based replication**.
  See Also master server, row-based replication, slave server, statement-based replication.

restore
: The process of putting a set of backup files from the **MySQL Enterprise Backup** product in place for use by MySQL. This operation can be performed to fix a corrupted database, to return to some earlier point in time, or (in a **replication** context) to set up a new **slave database**. In the **MySQL Enterprise Backup** product, this operation is performed by the `copy-back` option of the `mysqlbackup` command.
  See Also hot backup, MySQL Enterprise Backup, mysqlbackup command, prepared backup, replication, slave server.

rollback
: A **SQL** statement that ends a **transaction**, undoing any changes made by the transaction. It is the opposite of **commit**, which makes permanent any changes made in the transaction.

  By default, MySQL uses the **autocommit** setting, which automatically issues a commit following each SQL statement. You must change this setting before you can use the rollback technique.
  See Also ACID, autocommit, commit, SQL, transaction.

rollback segment
: The storage area containing the **undo log**. Rollback segments have traditionally resided in the **system tablespace**. As of MySQL 5.6, rollback segments can reside in separate **undo tablespaces**.
  See Also system tablespace, undo log, undo tablespace.

row
: The logical data structure defined by a set of **columns**. A set of rows makes up a **table**. Within `InnoDB` **data files**, each **page** can contain one or more rows.

  Although `InnoDB` uses the term **row format** for consistency with MySQL syntax, the row format is a property of each table and applies to all rows in that table.
  See Also column, data files, page, row format, table.

row format
: The disk storage format for **rows** of an `InnoDB` **table**. As `InnoDB` gains new capabilities such as **compression**, new row formats are introduced to support the resulting improvements in storage efficiency and performance.

  The row format of an `InnoDB` table is specified by the `ROW_FORMAT` option or by the `innodb_default_row_format` configuration option (introduced in MySQL 5.7.9). Row formats include `REDUNDANT`, `COMPACT`, `COMPRESSED`, and `DYNAMIC`. To view the row format of an `InnoDB` table, issue the `SHOW TABLE STATUS` statement or query `INFORMATION_SCHEMA.INNODB_SYS_TABLES` (available in MySQL 5.6 or higher).
  See Also compact row format, compressed row format, compression, dynamic row format, redundant row format, row, table.

row lock

A **lock** that prevents a row from being accessed in an incompatible way by another **transaction**. Other rows in the same table can be freely written to by other transactions. This is the type of **locking** done by **DML** operations on **InnoDB** tables.

Contrast with **table locks** used by `MyISAM`, or during **DDL** operations on `InnoDB` tables that cannot be done with **online DDL**; those locks block concurrent access to the table.
See Also DDL, DML, InnoDB, lock, locking, online DDL, table lock, transaction.

row-based replication

A form of **replication** where events are propagated from the **master server** specifying how to change individual rows on the **slave server**. It is safe to use for all settings of the `innodb_autoinc_lock_mode` option.
See Also auto-increment locking, innodb_autoinc_lock_mode, master server, replication, slave server, statement-based replication.

row-level locking

The **locking** mechanism used for **InnoDB** tables, relying on **row locks** rather than **table locks**. Multiple **transactions** can modify the same table concurrently. Only if two transactions try to modify the same row does one of the transactions wait for the other to complete (and release its row locks).
See Also InnoDB, locking, row lock, table lock, transaction.

rw-lock

The low-level object that `InnoDB` uses to represent and enforce shared-access **locks** to internal in-memory data structures following certain rules. Contrast with **mutexes**, which `InnoDB` uses to represent and enforce exclusive access to internal in-memory data structures. Mutexes and rw-locks are known collectively as **latches**.

`rw-lock` types include `s-locks` (shared locks), `x-locks` (exclusive locks), and `sx-locks` (shared-exclusive locks).

- An `s-lock` provides read access to a common resource.
- An `x-lock` provides write access to a common resource while not permitting inconsistent reads by other threads.
- An `sx-lock` provides write access to a common resource while permitting inconsistent reads by other threads. `sx-locks` were introduced in MySQL 5.7 to optimize concurrency and improve scalability for read-write workloads.

The following matrix summarizes rw-lock type compatibility.

	S	***SX***	***X***
S	Compatible	Compatible	Conflict
SX	Compatible	Conflict	Conflict
X	Conflict	Conflict	Conflict

See Also latch, lock, mutex, Performance Schema.

S

savepoint

Savepoints help to implement nested **transactions**. They can be used to provide scope to operations on tables that are part of a larger transaction. For example, scheduling a trip in a reservation system might involve booking several different flights; if a desired flight is unavailable, you might **roll back** the changes involved in booking that one leg, without rolling back the earlier flights that were successfully booked.
See Also rollback, transaction.

scalability
: The ability to add more work and issue more simultaneous requests to a system, without a sudden drop in performance due to exceeding the limits of system capacity. Software architecture, hardware configuration, application coding, and type of workload all play a part in scalability. When the system reaches its maximum capacity, popular techniques for increasing scalability are **scale up** (increasing the capacity of existing hardware or software) and **scale out** (adding new servers and more instances of MySQL). Often paired with **availability** as critical aspects of a large-scale deployment.
See Also availability, scale out, scale up.

scale out
: A technique for increasing **scalability** by adding new servers and more instances of MySQL. For example, setting up replication, NDB Cluster, connection pooling, or other features that spread work across a group of servers. Contrast with **scale up**.
See Also scalability, scale up.

scale up
: A technique for increasing **scalability** by increasing the capacity of existing hardware or software. For example, increasing the memory on a server and adjusting memory-related parameters such as `innodb_buffer_pool_size` and `innodb_buffer_pool_instances`. Contrast with **scale out**.
See Also scalability, scale out.

schema
: Conceptually, a schema is a set of interrelated database objects, such as tables, table columns, data types of the columns, indexes, foreign keys, and so on. These objects are connected through SQL syntax, because the columns make up the tables, the foreign keys refer to tables and columns, and so on. Ideally, they are also connected logically, working together as part of a unified application or flexible framework. For example, the **INFORMATION_SCHEMA** and **performance_schema** databases use “schema” in their names to emphasize the close relationships between the tables and columns they contain.

  In MySQL, physically, a **schema** is synonymous with a **database**. You can substitute the keyword `SCHEMA` instead of `DATABASE` in MySQL SQL syntax, for example using `CREATE SCHEMA` instead of `CREATE DATABASE`.

  Some other database products draw a distinction. For example, in the Oracle Database product, a **schema** represents only a part of a database: the tables and other objects owned by a single user.
  See Also database, INFORMATION_SCHEMA, Performance Schema.

SDI
: Acronym for “serialized dictionary information”.
See Also Serialized Dictionary Information (SDI).

search index
: In MySQL, **full-text search** queries use a special kind of index, the **FULLTEXT index**. In MySQL 5.6.4 and up, `InnoDB` and `MyISAM` tables both support `FULLTEXT` indexes; formerly, these indexes were only available for `MyISAM` tables.
See Also full-text search, FULLTEXT index.

secondary index
: A type of `InnoDB` **index** that represents a subset of table columns. An `InnoDB` table can have zero, one, or many secondary indexes. (Contrast with the **clustered index**, which is required for each `InnoDB` table, and stores the data for all the table columns.)

  A secondary index can be used to satisfy queries that only require values from the indexed columns. For more complex queries, it can be used to identify the relevant rows in the table, which are then retrieved through lookups using the clustered index.

  Creating and dropping secondary indexes has traditionally involved significant overhead from copying all the data in the `InnoDB` table. The **fast index creation** feature makes both `CREATE INDEX` and `DROP INDEX` statements much faster for `InnoDB` secondary indexes.
  See Also clustered index, Fast Index Creation, index.

segment
A division within an InnoDB **tablespace**. If a tablespace is analogous to a directory, the segments are analogous to files within that directory. A segment can grow. New segments can be created.

For example, within a **file-per-table** tablespace, table data is in one segment and each associated index is in its own segment. The **system tablespace** contains many different segments, because it can hold many tables and their associated indexes. The system tablespace also includes one or more **rollback segments** used for **undo logs**.

Segments grow and shrink as data is inserted and deleted. When a segment needs more room, it is extended by one **extent** (1 megabyte) at a time. Similarly, a segment releases one extent's worth of space when all the data in that extent is no longer needed.
See Also extent, file-per-table, rollback segment, system tablespace, tablespace, undo log.

selectivity
A property of data distribution, the number of distinct values in a column (its **cardinality**) divided by the number of records in the table. High selectivity means that the column values are relatively unique, and can retrieved efficiently through an index. If you (or the query optimizer) can predict that a test in a WHERE clause only matches a small number (or proportion) of rows in a table, the overall **query** tends to be efficient if it evaluates that test first, using an index.
See Also cardinality, query.

semi-consistent read
A type of read operation used for UPDATE statements, that is a combination of **READ COMMITTED** and **consistent read**. When an UPDATE statement examines a row that is already locked, InnoDB returns the latest committed version to MySQL so that MySQL can determine whether the row matches the WHERE condition of the UPDATE. If the row matches (must be updated), MySQL reads the row again, and this time InnoDB either locks it or waits for a lock on it. This type of read operation can only happen when the transaction has the READ COMMITTED **isolation level**, or when the innodb_locks_unsafe_for_binlog option is enabled. innodb_locks_unsafe_for_binlog was removed in MySQL 8.0.
See Also consistent read, isolation level, READ COMMITTED.

SERIALIZABLE
The **isolation level** that uses the most conservative locking strategy, to prevent any other **transactions** from inserting or changing data that was read by this transaction, until it is finished. This way, the same query can be run over and over within a transaction, and be certain to retrieve the same set of results each time. Any attempt to change data that was committed by another transaction since the start of the current transaction, cause the current transaction to wait.

This is the default isolation level specified by the SQL standard. In practice, this degree of strictness is rarely needed, so the default isolation level for InnoDB is the next most strict, **REPEATABLE READ**.
See Also ACID, consistent read, isolation level, locking, REPEATABLE READ, transaction.

Serialized Dictionary Information (SDI)
Dictionary object metadata in serialized BLOB form.
See Also file-per-table, general tablespace, system tablespace, tablespace.

server
A type of program that runs continuously, waiting to receive and act upon requests from another program (the **client**). Because often an entire computer is dedicated to running one or more server programs (such as a database server, a web server, an application server, or some combination of these), the term **server** can also refer to the computer that runs the server software.
See Also client, mysqld.

shared lock
A kind of **lock** that allows other **transactions** to read the locked object, and to also acquire other shared locks on it, but not to write to it. The opposite of **exclusive lock**.
See Also exclusive lock, lock, transaction.

shared tablespace
: Another way of referring to the **system tablespace** or to a **general tablespace**. General tablespaces were introduced in MySQL 5.7
See Also general tablespace, system tablespace.

sharp checkpoint
: The process of **flushing** to disk all **dirty** buffer pool pages whose redo entries are contained in certain portion of the **redo log**. Occurs before `InnoDB` reuses a portion of a log file; the log files are used in a circular fashion. Typically occurs with write-intensive **workloads**.
See Also dirty page, flush, redo log, workload.

shutdown
: The process of stopping the MySQL server. By default, this process cleans up operations for **InnoDB** tables, so `InnoDB` can be **slow** to shut down, but fast to start up later. If you skip the cleanup operations, it is **fast** to shut down but the cleanup must be performed during the next restart.

  The shutdown mode for `InnoDB` is controlled by the `innodb_fast_shutdown` option.
See Also fast shutdown, InnoDB, slow shutdown, startup.

slave server
: Frequently shortened to “slave”. A database **server** machine in a **replication** scenario that receives changes from another server (the **master**) and applies those same changes. Thus it maintains the same contents as the master, although it might lag somewhat behind.

  In MySQL, slave servers are commonly used in disaster recovery, to take the place of a master servers that fails. They are also commonly used for testing software upgrades and new settings, to ensure that database configuration changes do not cause problems with performance or reliability.

  Slave servers typically have high workloads, because they process all the **DML** (write) operations relayed from the master, as well as user queries. To ensure that slave servers can apply changes from the master fast enough, they frequently have fast I/O devices and sufficient CPU and memory to run multiple database instances on the same slave server. For example, the master server might use hard drive storage while the slave servers use **SSD**s.
See Also DML, master server, replication, server, SSD.

slow query log
: A type of **log** used for performance tuning of SQL statements processed by the MySQL server. The log information is stored in a file. You must enable this feature to use it. You control which categories of “slow” SQL statements are logged. For more information, see Section 5.4.5, “The Slow Query Log”.
See Also general query log, log.

slow shutdown
: A type of **shutdown** that does additional `InnoDB` flushing operations before completing. Also known as a **clean shutdown**. Specified by the configuration parameter `innodb_fast_shutdown=0` or the command `SET GLOBAL innodb_fast_shutdown=0;`. Although the shutdown itself can take longer, that time will be saved on the subsequent startup.
See Also clean shutdown, fast shutdown, shutdown.

snapshot
: A representation of data at a particular time, which remains the same even as changes are **committed** by other **transactions**. Used by certain **isolation levels** to allow **consistent reads**.
See Also commit, consistent read, isolation level, transaction.

sort buffer
: The buffer used for sorting data during creation of an `InnoDB` index. Sort buffer size is configured using the `innodb_sort_buffer_size` configuration option.

space ID
: An identifier used to uniquely identify an `InnoDB` **tablespace** within a MySQL instance. The space ID for the **system tablespace** is always zero; this same ID applies to all tables within the system tablespace or within a general tablespace. Each **file-per-table** tablespace and **general tablespace** has its own space ID.

Prior to MySQL 5.6, this hardcoded value presented difficulties in moving `InnoDB` tablespace files between MySQL instances. Starting in MySQL 5.6, you can copy tablespace files between instances by using the **transportable tablespace** feature involving the statements `FLUSH TABLES ... FOR EXPORT`, `ALTER TABLE ... DISCARD TABLESPACE`, and `ALTER TABLE ... IMPORT TABLESPACE`. The information needed to adjust the space ID is conveyed in the **.cfg file** which you copy along with the tablespace. See Section 15.7.6, "Copying File-Per-Table Tablespaces to Another Instance" for details.
See Also .cfg file, file-per-table, general tablespace, .ibd file, system tablespace, tablespace, transportable tablespace.

sparse file
A type of file that uses file system space more efficiently by writing metadata representing empty blocks to disk instead of writing the actual empty space. The `InnoDB` **transparent page compression** feature relies on sparse file support. For more information, see Section 15.9.2, "InnoDB Page Compression".
See Also hole punching, transparent page compression.

spin
A type of **wait** operation that continuously tests whether a resource becomes available. This technique is used for resources that are typically held only for brief periods, where it is more efficient to wait in a "busy loop" than to put the thread to sleep and perform a context switch. If the resource does not become available within a short time, the spin loop ceases and another wait technique is used.
See Also latch, lock, mutex, wait.

SQL
The Structured Query Language that is standard for performing database operations. Often divided into the categories **DDL**, **DML**, and **queries**. MySQL includes some additional statement categories such as **replication**. See Chapter 9, *Language Structure* for the building blocks of SQL syntax, Chapter 11, *Data Types* for the data types to use for MySQL table columns, Chapter 13, *SQL Statement Syntax* for details about SQL statements and their associated categories, and Chapter 12, *Functions and Operators* for standard and MySQL-specific functions to use in queries.
See Also DDL, DML, query, replication.

SSD
Acronym for "solid-state drive". A type of storage device with different performance characteristics than a traditional hard disk drive (**HDD**): smaller storage capacity, faster for random reads, no moving parts, and with a number of considerations affecting write performance. Its performance characteristics can influence the throughput of a **disk-bound** workload.
See Also disk-bound, HDD.

startup
The process of starting the MySQL server. Typically done by one of the programs listed in Section 4.3, "MySQL Server and Server-Startup Programs". The opposite of **shutdown**.
See Also shutdown.

statement-based replication
A form of **replication** where SQL statements are sent from the **master server** and replayed on the **slave server**. It requires some care with the setting for the `innodb_autoinc_lock_mode` option, to avoid potential timing problems with **auto-increment locking**.
See Also auto-increment locking, innodb_autoinc_lock_mode, master server, replication, row-based replication, slave server.

statistics
Estimated values relating to each `InnoDB` **table** and **index**, used to construct an efficient **query execution plan**. The main values are the **cardinality** (number of distinct values) and the total number of table rows or index entries. The statistics for the table represent the data in its **primary key** index. The statistics for a **secondary index** represent the rows covered by that index.

The values are estimated rather than counted precisely because at any moment, different **transactions** can be inserting and deleting rows from the same table. To keep the values from being recalculated frequently, you can enable **persistent statistics**, where the values are stored in `InnoDB` system tables, and refreshed only when you issue an `ANALYZE TABLE` statement.

You can control how **NULL** values are treated when calculating statistics through the `innodb_stats_method` configuration option.

Other types of statistics are available for database objects and database activity through the **INFORMATION_SCHEMA** and **PERFORMANCE_SCHEMA** tables.
See Also cardinality, index, INFORMATION_SCHEMA, NULL, Performance Schema, persistent statistics, primary key, query execution plan, secondary index, table, transaction.

stemming
The ability to search for different variations of a word based on a common root word, such as singular and plural, or past, present, and future verb tense. This feature is currently supported in `MyISAM` **full-text search** feature but not in **FULLTEXT indexes** for `InnoDB` tables.
See Also full-text search, FULLTEXT index.

stopword
In a **FULLTEXT index**, a word that is considered common or trivial enough that it is omitted from the **search index** and ignored in search queries. Different configuration settings control stopword processing for `InnoDB` and `MyISAM` tables. See Section 12.9.4, "Full-Text Stopwords" for details.
See Also FULLTEXT index, search index.

storage engine
A component of the MySQL database that performs the low-level work of storing, updating, and querying data. In MySQL 5.5 and higher, **InnoDB** is the default storage engine for new tables, superceding `MyISAM`. Different storage engines are designed with different tradeoffs between factors such as memory usage versus disk usage, read speed versus write speed, and speed versus robustness. Each storage engine manages specific tables, so we refer to `InnoDB` tables, `MyISAM` tables, and so on.

The **MySQL Enterprise Backup** product is optimized for backing up `InnoDB` tables. It can also back up tables handled by `MyISAM` and other storage engines.
See Also InnoDB, MySQL Enterprise Backup, table type.

stored generated column
A column whose values are computed from an expression included in the column definition. Column values are evaluated and stored when rows are inserted or updated. A stored generated column requires storage space and can be indexed.

Contrast with **virtual generated column**.
See Also base column, generated column, virtual generated column.

strict mode
The general name for the setting controlled by the `innodb_strict_mode` option. Turning on this setting causes certain conditions that are normally treated as warnings, to be considered errors. For example, certain invalid combinations of options related to **file format** and **row format**, that normally produce a warning and continue with default values, now cause the `CREATE TABLE` operation to fail. `innodb_strict_mode` is enabled by default in MySQL 5.7.

MySQL also has something called strict mode.
See Also file format, innodb_strict_mode, row format.

sublist
Within the list structure that represents the **buffer pool**, pages that are relatively old and relatively new are represented by different portions of the **list**. A set of parameters control the size of these portions and the dividing point between the new and old pages.
See Also buffer pool, eviction, list, LRU.

supremum record
A **pseudo-record** in an index, representing the **gap** above the largest value in that index. If a transaction has a statement such as `SELECT ... FROM ... WHERE col > 10 FOR UPDATE;`, and the largest value in the column is 20, it is a lock on the supremum record that prevents other transactions from inserting even larger values such as 50, 100, and so on.

See Also gap, infimum record, pseudo-record.

surrogate key
Synonym name for **synthetic key**.
See Also synthetic key.

synthetic key
An indexed column, typically a **primary key**, where the values are assigned arbitrarily. Often done using an **auto-increment** column. By treating the value as completely arbitrary, you can avoid overly restrictive rules and faulty application assumptions. For example, a numeric sequence representing employee numbers might have a gap if an employee was approved for hiring but never actually joined. Or employee number 100 might have a later hiring date than employee number 500, if they left the company and later rejoined. Numeric values also produce shorter values of predictable length. For example, storing numeric codes meaning “Road”, “Boulevard”, “Expressway”, and so on is more space-efficient than repeating those strings over and over.

Also known as a **surrogate key**. Contrast with **natural key**.
See Also auto-increment, natural key, primary key, surrogate key.

system tablespace
One or more data files (**ibdata files**) containing the metadata for `InnoDB`-related objects, and the storage areas for the **change buffer**, the **doublewrite buffer**, and possibly **undo logs**. It may also contain table and index data for `InnoDB` tables if tables were created in the system tablespace instead of **file-per-table** or **general tablespaces**. The data and metadata in the system tablespace apply to all **databases** in a MySQL **instance**.

Prior to MySQL 5.6.7, the default was to keep all `InnoDB` tables and indexes inside the system tablespace, often causing this file to become very large. Because the system tablespace never shrinks, storage problems could arise if large amounts of temporary data were loaded and then deleted. In MySQL 8.0, the default is **file-per-table** mode, where each table and its associated indexes are stored in a separate **.ibd file**. This new default makes it easier to use `InnoDB` features that rely on `DYNAMIC` and `COMPRESSED` row formats, such as table **compression**, efficient storage of **off-page columns**, and large index key prefixes.

The `innodb_undo_tablespaces` option allows you to configure separate tablespace files for undo logs.

Keeping all table data in the system tablespace or in separate `.ibd` files has implications for storage management in general. The **MySQL Enterprise Backup** product might back up a small set of large files, or many smaller files. On systems with thousands of tables, the file system operations to process thousands of `.ibd` files can cause bottlenecks.

`InnoDB` introduced general tablespaces in MySQL 5.7.6, which are also represented by `.ibd` files. General tablespaces are shared tablespaces created using `CREATE TABLESPACE` syntax. They can be created outside of the MySQL data directory, are capable of holding multiple tables, and support tables of all row formats.
See Also change buffer, compression, data dictionary, database, doublewrite buffer, dynamic row format, file-per-table, general tablespace, .ibd file, ibdata file, innodb_file_per_table, instance, MySQL Enterprise Backup, off-page column, tablespace, undo log.

T

table
Each MySQL table is associated with a particular **storage engine**. **InnoDB** tables have particular **physical** and **logical** characteristics that affect performance, **scalability**, **backup**, administration, and application development.

In terms of file storage, an `InnoDB` table belongs to one of the following tablespace types:

- The shared `InnoDB` **system tablespace**, which is comprised of one or more **ibdata files**.

- A **file-per-table** tablespace, comprised of an individual **.ibd file**.
- A shared **general tablespace**, comprised of an individual `.ibd` file. General tablespaces were introduced in MySQL 5.7.6.

`.ibd` data files contain both table and **index** data.

`InnoDB` tables created in file-per-table tablespaces can use **DYNAMIC** or **COMPRESSED** row format. These row formats enable `InnoDB` features such as **compression**, efficient storage of **off-page columns**, and large index key prefixes. General tablespaces support all row formats.

The system tablespace supports tables that use **REDUNDANT**, **COMPACT**, and **DYNAMIC** row formats. System tablespace support for the **DYNAMIC** row format was added in MySQL 5.7.6.

The **rows** of an `InnoDB` table are organized into an index structure known as the **clustered index**, with entries sorted based on the **primary key** columns of the table. Data access is optimized for queries that filter and sort on the primary key columns, and each index contains a copy of the associated primary key columns for each entry. Modifying values for any of the primary key columns is an expensive operation. Thus an important aspect of `InnoDB` table design is choosing a primary key with columns that are used in the most important queries, and keeping the primary key short, with rarely changing values.
See Also backup, clustered index, compact row format, compressed row format, compression, dynamic row format, Fast Index Creation, file-per-table, .ibd file, index, off-page column, primary key, redundant row format, row, system tablespace, tablespace.

table lock
: A lock that prevents any other **transaction** from accessing a table. `InnoDB` makes considerable effort to make such locks unnecessary, by using techniques such as **online DDL**, **row locks** and **consistent reads** for processing **DML** statements and **queries**. You can create such a lock through SQL using the `LOCK TABLE` statement; one of the steps in migrating from other database systems or MySQL storage engines is to remove such statements wherever practical.
See Also consistent read, DML, lock, locking, online DDL, query, row lock, table, transaction.

table scan
: See full table scan.

table statistics
: See statistics.

table type
: Obsolete synonym for **storage engine**. We refer to `InnoDB` tables, `MyISAM` tables, and so on.
See Also InnoDB, storage engine.

tablespace
: A data file that can hold data for one or more `InnoDB` **tables** and associated **indexes**.

The **system tablespace** contains the `InnoDB` **data dictionary**, and prior to MySQL 5.6 holds all other `InnoDB` tables by default.

The `innodb_file_per_table` option, enabled by default in MySQL 5.6 and higher, allows tables to be created in their own tablespaces. File-per-table tablespaces support features such as efficient storage of **off-page columns**, table compression, and transportable tablespaces. See Section 15.7.4, “InnoDB File-Per-Table Tablespaces” for details.

`InnoDB` introduced general tablespaces in MySQL 5.7.6. General tablespaces are shared tablespaces created using `CREATE TABLESPACE` syntax. They can be created outside of the MySQL data directory, are capable of holding multiple tables, and support tables of all row formats.

MySQL NDB Cluster also groups its tables into tablespaces. See NDB Cluster Disk Data Objects for details.
See Also compressed row format, data dictionary, data files, file-per-table, general tablespace, index, innodb_file_per_table, system tablespace, table.

tablespace dictionary
: A tablespace dictionary present in persistent `InnoDB` tablespaces. It stores Serialized Dictionary Information (SDI) for objects contained within the tablespace.
See Also data dictionary, file-per-table, .frm file, general tablespace, .ibd file, system tablespace, tablespace.

temporary table
: A **table** whose data does not need to be truly permanent. For example, temporary tables might be used as storage areas for intermediate results in complicated calculations or transformations; this intermediate data would not need to be recovered after a crash. Database products can take various shortcuts to improve the performance of operations on temporary tables, by being less scrupulous about writing data to disk and other measures to protect the data across restarts.

Sometimes, the data itself is removed automatically at a set time, such as when the transaction ends or when the session ends. With some database products, the table itself is removed automatically too.
See Also table.

temporary tablespace
: The tablespace for non-compressed `InnoDB` **temporary tables** and related objects, introduced in MySQL 5.7. The `innodb_temp_data_file_path` configuration file option defines a relative path for the temporary tablespace data file. If `innodb_temp_data_file_path` is not specified, the default behavior is to create a single auto-extending 12MB data file named `ibtmp1` in the data directory. The temporary tablespace is recreated on each server start and receives a dynamically generated **space ID**, which helps avoid conflicts with existing space IDs. The temporary tablespace cannot reside on a raw device. Startup is refused if the temporary tablespace cannot be created.

The temporary tablespace is removed on normal shutdown or on an aborted initialization. The temporary tablespace is not removed when a crash occurs. In this case, the database administrator may remove the temporary tablespace manually or restart the server with the same configuration, which removes and recreates the temporary tablespace.
See Also ibtmp file, space ID, system tablespace, temporary table.

text collection
: The set of columns included in a **FULLTEXT index**.
See Also FULLTEXT index.

thread
: A unit of processing that is typically more lightweight than a **process**, allowing for greater **concurrency**.
See Also concurrency, master thread, process, Pthreads.

torn page
: An error condition that can occur due to a combination of I/O device configuration and hardware failure. If data is written out in chunks smaller than the `InnoDB` **page size** (by default, 16KB), a hardware failure while writing could result in only part of a page being stored to disk. The `InnoDB` **doublewrite buffer** guards against this possibility.
See Also doublewrite buffer.

TPS
: Acronym for “**transactions** per second”, a unit of measurement sometimes used in benchmarks. Its value depends on the **workload** represented by a particular benchmark test, combined with factors that you control such as the hardware capacity and database configuration.
See Also transaction, workload.

transaction
: Transactions are atomic units of work that can be **committed** or **rolled back**. When a transaction makes multiple changes to the database, either all the changes succeed when the transaction is committed, or all the changes are undone when the transaction is rolled back.

Database transactions, as implemented by `InnoDB`, have properties that are collectively known by the acronym **ACID**, for atomicity, consistency, isolation, and durability.
See Also ACID, commit, isolation level, lock, rollback.

transaction ID
An internal field associated with each **row**. This field is physically changed by `INSERT`, `UPDATE`, and `DELETE` operations to record which **transaction** has locked the row.
See Also implicit row lock, row, transaction.

transparent page compression
A feature added in MySQL 5.7.8 that permits page-level compression for `InnoDB` tables that reside in **file-per-table** tablespaces. Page compression is enabled by specifying the `COMPRESSION` attribute with `CREATE TABLE` or `ALTER TABLE`. For more information, see Section 15.9.2, “InnoDB Page Compression”.
See Also file-per-table, hole punching, sparse file.

transportable tablespace
A feature that allows a **tablespace** to be moved from one instance to another. Traditionally, this has not been possible for `InnoDB` tablespaces because all table data was part of the **system tablespace**. In MySQL 5.6 and higher, the `FLUSH TABLES ... FOR EXPORT` syntax prepares an `InnoDB` table for copying to another server; running `ALTER TABLE ... DISCARD TABLESPACE` and `ALTER TABLE ... IMPORT TABLESPACE` on the other server brings the copied data file into the other instance. A separate **.cfg file**, copied along with the **.ibd file**, is used to update the table metadata (for example the **space ID**) as the tablespace is imported. See Section 15.7.6, “Copying File-Per-Table Tablespaces to Another Instance” for usage information.
See Also .cfg file, .ibd file, space ID, system tablespace, tablespace.

troubleshooting
The process of determining the source of a problem. Some of the resources for troubleshooting MySQL problems include:

- Section 2.9.2.1, “Troubleshooting Problems Starting the MySQL Server”
- Section 6.2.9, “Troubleshooting Problems Connecting to MySQL”
- Section B.5.3.2, “How to Reset the Root Password”
- Section B.5.2, “Common Errors When Using MySQL Programs”
- Section 15.20, “InnoDB Troubleshooting”.

truncate
A **DDL** operation that removes the entire contents of a table, while leaving the table and related indexes intact. Contrast with **drop**. Although conceptually it has the same result as a `DELETE` statement with no `WHERE` clause, it operates differently behind the scenes: `InnoDB` creates a new empty table, drops the old table, then renames the new table to take the place of the old one. Because this is a DDL operation, it cannot be **rolled back**.

If the table being truncated contains **foreign keys** that reference another table, the truncation operation uses a slower method of operation, deleting one row at a time so that corresponding rows in the referenced table can be deleted as needed by any `ON DELETE CASCADE` clause. (MySQL 5.5 and higher do not allow this slower form of truncate, and return an error instead if foreign keys are involved. In this case, use a `DELETE` statement instead.
See Also DDL, drop, foreign key, rollback.

tuple
A technical term designating an ordered set of elements. It is an abstract notion, used in formal discussions of database theory. In the database field, tuples are usually represented by the columns of a table row. They could also be represented by the result sets of queries, for example, queries that retrieved only some columns of a table, or columns from joined tables.
See Also cursor.

two-phase commit
An operation that is part of a distributed **transaction**, under the **XA** specification. (Sometimes abbreviated as 2PC.) When multiple databases participate in the transaction, either all databases **commit** the changes, or all databases **roll back** the changes.

See Also commit, rollback, transaction, XA.

U

undo
: Data that is maintained throughout the life of a **transaction**, recording all changes so that they can be undone in case of a **rollback** operation. It is stored in the **undo log** either within the **system tablespace** or in separate **undo tablespaces**.
See Also rollback, rollback segment, system tablespace, transaction, undo log, undo tablespace.

undo buffer
: See undo log.

undo log
: A storage area that holds copies of data modified by active **transactions**. If another transaction needs to see the original data (as part of a **consistent read** operation), the unmodified data is retrieved from this storage area.

In MySQL 5.6 and higher, you can use the `innodb_undo_tablespaces` configuration option to move undo logs away from the system tablespace to one or more separate **undo tablespace** files, optionally stored on another storage device such as an **SSD**.

The undo log is split into separate portions, the **insert undo buffer** and the **update undo buffer**.
See Also consistent read, rollback segment, SSD, system tablespace, transaction, undo tablespace.

undo log segment
: A collection of **undo logs**. Undo log segments exists within **rollback segments**. An undo log segment might contain undo logs from multiple transactions. An undo log segment can only be used by one transaction at a time but can be reused after it is released at transaction **commit** or **rollback**. May also be referred to as an "undo segment".
See Also commit, rollback, rollback segment, undo log.

undo tablespace
: An undo tablespace comprises one or more files that contain **undo logs**. Undo logs exist within **undo log segments**, which are contained within **rollback segments**. Rollback segments have traditionally resided in the system tablespace. As of MySQL 5.6, rollback segments can reside in separate undo tablespaces. An undo tablespace is created when undo logs are separated from the system tablespace using the `innodb_undo_tablespaces` configuration option.

For more information, see Section 15.7.7, "Storing InnoDB Undo Logs in Separate Tablespaces".
See Also rollback segment, system tablespace, undo log, undo log segment.

unique constraint
: A kind of **constraint** that asserts that a column cannot contain any duplicate values. In terms of **relational** algebra, it is used to specify 1-to-1 relationships. For efficiency in checking whether a value can be inserted (that is, the value does not already exist in the column), a unique constraint is supported by an underlying **unique index**.
See Also constraint, relational, unique index.

unique index
: An index on a column or set of columns that have a **unique constraint**. Because the index is known not to contain any duplicate values, certain kinds of lookups and count operations are more efficient than in the normal kind of index. Most of the lookups against this type of index are simply to determine if a certain value exists or not. The number of values in the index is the same as the number of rows in the table, or at least the number of rows with non-null values for the associated columns.

Change buffering optimization does not apply to unique indexes. As a workaround, you can temporarily set `unique_checks=0` while doing a bulk data load into an `InnoDB` table.
See Also cardinality, change buffering, unique constraint, unique key.

unique key
The set of columns (one or more) comprising a **unique index**. When you can define a `WHERE` condition that matches exactly one row, and the query can use an associated unique index, the lookup and error handling can be performed very efficiently.
See Also cardinality, unique constraint, unique index.

V

variable-length type
A data type of variable length. `VARCHAR`, `VARBINARY`, and `BLOB` and `TEXT` types are variable-length types.

`InnoDB` treats fixed-length fields greater than or equal to 768 bytes in length as variable-length fields, which can be stored **off-page**. For example, a `CHAR(255)` column can exceed 768 bytes if the maximum byte length of the character set is greater than 3, as it is with `utf8mb4`.
See Also off-page column, overflow page.

victim
The **transaction** that is automatically chosen to be **rolled back** when a **deadlock** is detected. `InnoDB` rolls back the transaction that has updated the fewest rows.

Deadlock detection can be disabled using the `innodb_deadlock_detect` configuration option.
See Also deadlock, deadlock detection, innodb_lock_wait_timeout, transaction.

virtual column
See virtual generated column.

virtual generated column
A column whose values are computed from an expression included in the column definition. Column values are not stored, but are evaluated when rows are read, immediately after any `BEFORE` triggers. A virtual generated column takes no storage. `InnoDB` supports secondary indexes on virtual generated columns.

Contrast with **stored generated column**.
See Also base column, generated column, stored generated column.

virtual index
A virtual index is a **secondary index** on one or more **virtual generated columns** or on a combination of virtual generated columns and regular columns or stored generated columns. For more information, see Section 13.1.15.9, “Secondary Indexes and Generated Columns”.
See Also secondary index, stored generated column, virtual generated column.

W

wait
When an operation, such as acquiring a **lock**, **mutex**, or **latch**, cannot be completed immediately, `InnoDB` pauses and tries again. The mechanism for pausing is elaborate enough that this operation has its own name, the **wait**. Individual threads are paused using a combination of internal `InnoDB` scheduling, operating system `wait()` calls, and short-duration **spin** loops.

On systems with heavy load and many transactions, you might use the output from the `SHOW INNODB STATUS` command or **Performance Schema** to determine whether threads are spending too much time waiting, and if so, how you can improve **concurrency**.
See Also concurrency, latch, lock, mutex, Performance Schema, spin.

warm backup
A **backup** taken while the database is running, but that restricts some database operations during the backup process. For example, tables might become read-only. For busy applications and web sites, you might prefer a **hot backup**.
See Also backup, cold backup, hot backup.

warm up
To run a system under a typical **workload** for some time after startup, so that the **buffer pool** and other memory regions are filled as they would be under normal conditions. This process happens naturally over time when a MySQL server is restarted or subjected to a new workload.

Typically, you run a workload for some time to warm up the buffer pool before running performance tests, to ensure consistent results across multiple runs; otherwise, performance might be artificially low during the first run.

In MySQL 5.6, you can speed up the warmup process by enabling the `innodb_buffer_pool_dump_at_shutdown` and `innodb_buffer_pool_load_at_startup` configuration options, to bring the contents of the buffer pool back into memory after a restart. These options are enabled by default in MySQL 5.7. See Section 15.6.3.8, “Saving and Restoring the Buffer Pool State”.
See Also buffer pool, workload.

workload
The combination and volume of **SQL** and other database operations, performed by a database application during typical or peak usage. You can subject the database to a particular workload during performance testing to identify **bottlenecks**, or during capacity planning.
See Also bottleneck, CPU-bound, disk-bound, SQL.

write combining
An optimization technique that reduces write operations when **dirty pages** are **flushed** from the `InnoDB` **buffer pool**. If a row in a page is updated multiple times, or multiple rows on the same page are updated, all of those changes are stored to the data files in a single write operation rather than one write for each change.
See Also buffer pool, dirty page, flush.

X

XA
A standard interface for coordinating distributed **transactions**, allowing multiple databases to participate in a transaction while maintaining **ACID** compliance. For full details, see Section 13.3.7, “XA Transactions”.

XA Distributed Transaction support is enabled by default. If you are not using this feature, you can disable the `innodb_support_xa` configuration option, avoiding the performance overhead of an extra fsync for each transaction.

As of MySQL 5.7.10, disabling `innodb_support_xa` is not permitted as it makes replication unsafe and prevents performance gains associated with **binary log** group commit. The `innodb_support_xa` configuration option is removed in MySQL 8.0.
See Also ACID, binary log, commit, transaction, two-phase commit.

Y

young
A characteristic of a **page** in the `InnoDB` **buffer pool** meaning it has been accessed recently, and so is moved within the buffer pool data structure, so that it will not be **flushed** soon by the **LRU** algorithm. This term is used in some **INFORMATION_SCHEMA** column names of tables related to the buffer pool.
See Also buffer pool, flush, INFORMATION_SCHEMA, LRU, page.